

# Young's Inequality for Increasing Functions

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## Abstract

Young's inequality states that

$$ab \leq \int_0^a f(x)dx + \int_0^b f^{-1}(y)dy$$

where  $a \geq 0, b \geq 0$  and  $f$  is strictly increasing and continuous. Its proof is formalised following the development by Cunningham and Grossman [1]. Their idea is to make the intuitive, geometric folklore proof rigorous by reasoning about step functions. The lack of the Riemann integral makes the development longer than one would like, but their argument is reproduced faithfully.

## Contents

<b>1 Young's Inequality for Increasing Functions</b>	<b>3</b>
1.1 Toward Young's inequality . . . . .	3
1.2 Regular divisions . . . . .	4
1.3 Special cases of Young's inequality . . . . .	7

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# 1 Young's Inequality for Increasing Functions

From the following paper: Cunningham, F., and Nathaniel Grossman. “On Young’s Inequality.” The American Mathematical Monthly 78, no. 7 (1971): 781–83. <https://doi.org/10.2307/2318018>

```
theory Youngs imports  
HOL-Analysis.Analysis
```

```
begin
```

## 1.1 Toward Young's inequality

```
lemma strict-mono-image-endpoints:  
  fixes f :: 'a::linear-continuum-topology ⇒ 'b::linorder-topology  
  assumes strict-mono-on {a..b} f and f: continuous-on {a..b} f and a ≤ b  
  shows f ` {a..b} = {f a..f b}  
proof  
  show f ` {a..b} ⊆ {f a..f b}  
    using assms(1) strict-mono-on-leD by fastforce  
  show {f a..f b} ⊆ f ` {a..b}  
    using assms IVT'[OF _ _ _ f] by (force simp: Bex-def)  
qed
```

Generalisations of the type of  $f$  are not obvious

```
lemma strict-mono-continuous-invD:  
  fixes f :: real ⇒ real  
  assumes sm: strict-mono-on {a..} f and contf: continuous-on {a..} f  
        and fim: f ` {a..} = {f a..} and g: ∀x. x ≥ a ⇒ g (f x) = x  
  shows continuous-on {f a..} g  
proof (clarify simp add: continuous-on-eq-continuous-within)  
  fix y  
  assume f a ≤ y  
  then obtain u where u: y+1 = f u u ≥ a  
    by (smt (verit, best) atLeast-iff fim imageE)  
  have continuous-on {f a..y+1} g  
  proof –  
    obtain continuous-on {a..u} f strict-mono-on {a..u} f  
      using contf sm continuous-on-subset by (force simp add: strict-mono-on-def)  
    moreover have continuous-on (f ` {a..u}) g  
      using assms continuous-on-subset  
      by (intro continuous-on-inv) fastforce+  
    ultimately show ?thesis  
      using strict-mono-image-endpoints [of _ _ f]  
      by (simp add: strict-mono-image-endpoints u)  
  qed  
  then have *: continuous (at y within {f a..y+1}) g  
    by (simp add: ‹f a ≤ y› continuous-on-imp-continuous-within)  
  show continuous (at y within {f a..}) g  
  proof (clarify simp add: continuous-within-topological Ball-def)
```

```

fix B
assume open B and g y ∈ B
with * obtain A where A: open A y ∈ A and ⋀x. f a ≤ x and x ≤ y+1 ⇒ x
    ∈ A → g x ∈ B
        by (force simp: continuous-within-topological)
    then have ∀x≥f a. x ∈ A ∩ ball y 1 → g x ∈ B
        by (smt (verit, ccfv-threshold) IntE dist-norm mem-ball real-norm-def)
    then show ∃A. open A and y ∈ A and (∀x≥f a. x ∈ A → g x ∈ B)
        by (metis Elementary-Metric-Spaces.open-ball Int-iff A centre-in-ball open-Int
            zero-less-one)
    qed
qed

```

## 1.2 Regular divisions

Our lack of the Riemann integral forces us to construct explicitly the step functions mentioned in the text.

**definition** segment ≡ λn k. {real k / real n..(1 + k) / real n}

**lemma** segment-nonempty: segment n k ≠ {}  
**by** (auto simp: segment-def divide-simps)

**lemma** segment-Suc: segment n ‘ {.. $<$ Suc k} = insert {k/n..(1 + real k) / n}  
 (segment n ‘ {.. $<$ k})  
**by** (simp add: segment-def lessThan-Suc)

**lemma** Union-segment-image: ⋃ (segment n ‘ {.. $<$ k}) = (if k=0 then {} else  
 {0..real k/real n})  
**proof** (induction k)  
**case** (Suc k)  
**then show** ?case  
**by** (simp add: divide-simps segment-Suc Un-commute ivl-disj-un-two-touch split:  
 if-split-asm)  
**qed** (auto simp: segment-def)

**definition** segments ≡ λn. segment n ‘ {.. $n$ }

**lemma** card-segments [simp]: card (segments n) = n  
**by** (simp add: segments-def segment-def card-image divide-right-mono inj-on-def)

**lemma** segments-0 [simp]: segments 0 = {}  
**by** (simp add: segments-def)

**lemma** Union-segments: ⋃ (segments n) = (if n=0 then {} else {0..1})  
**by** (simp add: segments-def Union-segment-image)

**definition** regular-division ≡ λa b n. (image ((+) a o (\*)(b-a))) ‘ (segments n)

**lemma** translate-scale-01:

```

assumes a ≤ b
shows (λx. a + (b - a) * x) ` {0..1} = {a..b::real}
using closed-segment-real-eq [of a b] assms closed-segment-eq-real-ivl by auto

lemma finite-regular-division [simp]: finite (regular-division a b n)
by (simp add: regular-division-def segments-def)

lemma card-regular-division [simp]:
assumes a < b
shows card (regular-division a b n) = n
proof -
have inj-on ((+) ((+ a) o (* (b - a))) (segments n))
proof
fix x y
assume ((+) a o (* (b - a))) ` x = ((+) a o (* (b - a))) ` y
then have (+ (-a)) ` ((+) a o (* (b - a))) ` x = (+ (-a)) ` ((+) a o (* (b - a))) ` y
by simp
then have ((* (b - a))) ` x = ((* (b - a))) ` y
by (simp add: image-comp)
then have (* (inverse(b - a))) ` (*) (b - a) ` x = (* (inverse(b - a))) ` (*) (b - a) ` y
by simp
then show x = y
using assms by (simp add: image-comp mult-ac)
qed
then show ?thesis
by (metis card-image card-segments regular-division-def)
qed

lemma Union-regular-division:
assumes a ≤ b
shows ∪(regular-division a b n) = (if n=0 then {} else {a..b})
using assms
by (auto simp: regular-division-def Union-segments translate-scale-01 simp flip: image-Union)

lemma regular-division-eqI:
assumes K: K = {a + (b-a)*(real k / n) .. a + (b-a)*((1 + real k) / n)}
and a < b k < n
shows K ∈ regular-division a b n
unfolding regular-division-def segments-def image-comp
proof
have K = (λx. (b-a) * x + a) ` {real k / real n..(1 + real k) / real n}
using K ` a < b by (simp add: image-affinity-atLeastAtMost divide-simps)
then show K = ((+) ((+ a) o (* (b - a))) o segment n) k
by (simp add: segment-def add.commute)
qed (use assms in auto)

```

```

lemma regular-divisionE:
  assumes K ∈ regular-division a b n a < b
  obtains k where k < n K = {a + (b-a)*(real k / n) .. a + (b-a)*((1 + real k)
  / n)}
  proof -
    have eq: (λx. a + (b - a) * x) = (λx. a + x) ∘ (λx. (b - a) * x)
      by (simp add: o-def)
    obtain k where k < n K = ((λx. a+x) ∘ (λx. (b-a) * x)) ` {k/n .. (1 + real k)
    / n}
      using assms by (auto simp: regular-division-def segments-def segment-def)
      with that ‹a < b› show ?thesis
        unfolding image-comp [symmetric] by auto
  qed

lemma regular-division-division-of:
  assumes a < b n > 0
  shows (regular-division a b n) division-of {a..b}
  proof (rule division-ofI)
    show finite (regular-division a b n)
      by (simp add: regular-division-def segments-def)
    show §: ⋃ (regular-division a b n) = {a..b}
      using Union-regular-division assms by simp
    fix K
    assume K: K ∈ regular-division a b n
    then obtain k where Keq: K = {a + (b-a)*(k/n) .. a + (b-a)*((1 + real k)
    / n)}
      using ‹a < b› regular-divisionE by meson
    show K ⊆ {a..b}
      using K Union-regular-division ‹n > 0› by (metis Union-upper §)
    show K ≠ {}
      using K by (auto simp: regular-division-def segment-nonempty segments-def)
    show ∃ a b. K = cbox a b
      by (metis K ‹a < b› box-real(2) regular-divisionE)
    fix K'
    assume K': K' ∈ regular-division a b n and K ≠ K'
    then obtain k' where Keq': K' = {a + (b-a)*(k'/n) .. a + (b-a)*((1 + real
    k') / n)}
      using K ‹a < b› regular-divisionE by meson
    consider 1 + real k ≤ k' | 1 + real k' ≤ k
      using Keq Keq' ‹K ≠ K'› by force
    then show interior K ∩ interior K' = {}
    proof cases
      case 1
      then show ?thesis
        by (simp add: Keq Keq' min-def max-def divide-right-mono assms)
      next
      case 2
      then have interior K' ∩ interior K = {}
        by (simp add: Keq Keq' min-def max-def divide-right-mono assms)
    qed
  qed

```

```

then show ?thesis
  by (simp add: inf-commute)
qed
qed

```

### 1.3 Special cases of Young's inequality

```

lemma weighted-nesting-sum:
  fixes g :: nat  $\Rightarrow$  'a::comm-ring-1
  shows  $(\sum k < n. (1 + of\text{-}nat k) * (g (Suc k) - g k)) = of\text{-}nat n * g n - (\sum i < n.$ 
 $g i)$ 
  by (induction n) (auto simp: algebra-simps)

theorem Youngs-exact:
  fixes f :: real  $\Rightarrow$  real
  assumes sm: strict-mono-on {0..} f and cont: continuous-on {0..} f and a:
 $a \geq 0$ 
  and f 0 = 0 f a = b
  and g:  $\bigwedge x. [0 \leq x; x \leq a] \implies g(f x) = x$ 
  shows a*b = integral {0..a} f + integral {0..b} g
  proof (cases a=0)
    case False
    with ‹a ≥ 0› have a > 0 by linarith
    then have b ≥ 0
      by (smt (verit, best) atLeast-iff sm strict-mono-onD)
    have sm-0a: strict-mono-on {0..a} f
      by (metis atLeastAtMost-iff atLeast-iff sm strict-mono-on-def)
    have cont-0a: continuous-on {0..a} f
      using cont continuous-on-subset by fastforce
    with sm-0a have continuous-on {0..b} g
      by (metis a atLeastAtMost-iff compact-Icc continuous-on-inv f g strict-mono-image-endpoints)
    then have intgb-g: g integrable-on {0..b}
      using integrable-continuous-interval by blast
    have intgb-f: f integrable-on {0..a}
      using cont-0a integrable-continuous-real by blast

    have f-iff [simp]: f x < f y  $\longleftrightarrow$  x < y f x ≤ f y  $\longleftrightarrow$  x ≤ y
    if x ≥ 0 y ≥ 0 for x y
      using that by (smt (verit, best) atLeast-iff sm strict-mono-onD) +
    have fim: f ` {0..a} = {0..b}
      by (simp add: ‹a ≥ 0› cont-0a strict-mono-image-endpoints strict-mono-on-def
f)
    have uniformly-continuous-on {0..a} f
      using compact-uniformly-continuous cont-0a by blast
    then obtain del where del-gt0:  $\bigwedge e. e > 0 \implies del e > 0$ 
      and del:  $\bigwedge e x x'. [|x' - x| < del e; e > 0; x \in \{0..a\}; x' \in \{0..a\}] \implies |f x' - f x| < e$ 
      unfolding uniformly-continuous-on-def dist-real-def by metis

```

```

have *:  $|a * b - \text{integral } \{0..a\} f - \text{integral } \{0..b\} g| < 2*\varepsilon$  if  $\varepsilon > 0$  for  $\varepsilon$ 
proof -
  define  $\delta$  where  $\delta = \min a (\text{del } (\varepsilon/a)) / 2$ 
  have  $\delta > 0 \ \delta \leq a$ 
    using  $\langle a > 0 \ \langle \varepsilon > 0 \rangle \ \text{del-gt0}$  by (auto simp:  $\delta$ -def)
  define  $n$  where  $n \equiv \text{nat}\lfloor a / \delta \rfloor$ 
  define  $a\text{-seg}$  where  $a\text{-seg} \equiv \lambda u::\text{real}. u * a/n$ 
  have  $n > 0$ 
    using  $\langle a > 0 \ \langle \delta > 0 \rangle \ \langle \delta \leq a \rangle$  by (simp add:  $n$ -def)
  have  $a\text{-seg-ge-0}$  [simp]:  $a\text{-seg } x \geq 0 \longleftrightarrow x \geq 0$ 
    and  $a\text{-seg-le-a}$  [simp]:  $a\text{-seg } x \leq a \longleftrightarrow x \leq n$  for  $x$ 
    using  $\langle n > 0 \ \langle a > 0 \rangle$  by (auto simp:  $a\text{-seg-def zero-le-mult-iff divide-simps}$ )
  have  $a\text{-seg-le-iff}$  [simp]:  $a\text{-seg } x \leq a\text{-seg } y \longleftrightarrow x \leq y$ 
    and  $a\text{-seg-less-iff}$  [simp]:  $a\text{-seg } x < a\text{-seg } y \longleftrightarrow x < y$  for  $x y$ 
    using  $\langle n > 0 \ \langle a > 0 \rangle$  by (auto simp:  $a\text{-seg-def zero-le-mult-iff divide-simps}$ )
  have strict-mono  $a\text{-seg}$ 
    by (simp add: strict-mono-def)
  have  $a\text{-seg-eq-a-iff}$ :  $a\text{-seg } x = a \longleftrightarrow x = n$  for  $x$ 
    using  $\langle 0 < n \ \langle a > 0 \rangle$  by (simp add:  $a\text{-seg-def nonzero-divide-eq-eq}$ )
  have  $f\text{a-eq-b}$ :  $f (a\text{-seg } n) = b$ 
    using  $a\text{-seg-eq-a-iff } f$  by fastforce

  have  $a/d < \text{real-of-int } \lfloor a * 2 / \min a d \rfloor$  if  $d > 0$  for  $d$ 
    by (smt (verit)  $\langle 0 < \delta \rangle \ \langle \delta \leq a \rangle$  add-divide-distrib divide-less-eq-1-pos floor-eq-iff
that)
  then have  $a\text{n-less-del}$ :  $a/n < \text{del } (\varepsilon/a)$ 
    using  $\langle a > 0 \ \langle \varepsilon > 0 \rangle \ \text{del-gt0}$  by (simp add:  $n$ -def  $\delta$ -def field-simps)

define  $lower$  where  $lower \equiv \lambda x. a\text{-seg}\lfloor (real n * x) / a \rfloor$ 
define  $f1$  where  $f1 \equiv f \circ lower$ 
have  $f1\text{-lower}$ :  $f1 x \leq f x$  if  $0 \leq x x \leq a$  for  $x$ 
proof -
  have  $lower x \leq x$ 
    using  $\langle n > 0 \rangle \ \text{floor-divide-lower}$  [OF  $\langle a > 0 \rangle$ ]
    by (auto simp: lower-def  $a\text{-seg-def field-simps}$ )
  moreover have  $lower x \geq 0$ 
    unfolding lower-def using  $\langle n > 0 \ \langle a \geq 0 \rangle \ \langle 0 \leq x \rangle$  by force
  ultimately show ?thesis
    using sm strict-mono-on-leD by (fastforce simp add:  $f1$ -def)
qed
define  $upper$  where  $upper \equiv \lambda x. a\text{-seg}\lceil real n * x / a \rceil$ 
define  $f2$  where  $f2 \equiv f \circ upper$ 
have  $f2\text{-upper}$ :  $f2 x \geq f x$  if  $0 \leq x x \leq a$  for  $x$ 
proof -
  have  $x \leq upper x$ 
    using  $\langle n > 0 \rangle \ \text{ceiling-divide-upper}$  [OF  $\langle a > 0 \rangle$ ] by (simp add: upper-def
 $a\text{-seg-def field-simps}$ )
  then show ?thesis
    using sm strict-mono-on-leD  $\langle 0 \leq x \rangle$  by (force simp:  $f2$ -def)

```

```

qed
let ?D = regular-division 0 a n
have div: ?D division-of {0..a}
  using ⟨a > 0⟩ ⟨n > 0⟩ regular-division-division-of zero-less-nat-eq by pres-
burger

have int-f1-D: (f1 has-integral f(Inf K) * (a/n)) K
  and int-f2-D: (f2 has-integral f(Sup K) * (a/n)) K and less: |f(Sup K) −
f(Inf K)| < ε/a
    if K ∈ ?D for K
proof -
  from regular-divisionE [OF that] ⟨a > 0⟩
  obtain k where k < n and k: K = {a-seg(real k)..a-seg(Suc k)}
    by (auto simp: a-seg-def mult.commute)
  define u where u ≡ a-seg k
  define v where v ≡ a-seg (Suc k)
  have u < v 0 ≤ u 0 ≤ v u ≤ a v ≤ a and Kuv: K = {u..v}
    using ⟨n > 0⟩ ⟨k < n⟩ ⟨a > 0⟩ by (auto simp: k u-def v-def divide-simps)
  have InfK: Inf K = u and SupK: Sup K = v
    using Kuv ⟨u < v⟩ apply force
    using ⟨n > 0⟩ ⟨a > 0⟩ by (auto simp: divide-right-mono k u-def v-def)
  have f1: f1 x = f (Inf K) if x ∈ K − {v} for x
  proof -
    have x ∈ {u..v}
      using that Kuv atLeastLessThan-eq-atLeastAtMost-diff by blast
      then have [real-of-int n * x / a] = int k
        using ⟨n > 0⟩ ⟨a > 0⟩ by (simp add: field-simps u-def v-def a-seg-def
floor-eq-iff)
      then show ?thesis
        by (simp add: InfK f1-def lower-def a-seg-def mult.commute u-def)
  qed
  have ((λx. f (Inf K)) has-integral (f (Inf K) * (a/n))) K
    using has-integral-const-real [of f (Inf K) u v]
      ⟨n > 0⟩ ⟨a > 0⟩ by (simp add: Kuv field-simps a-seg-def u-def v-def)
  then show (f1 has-integral (f (Inf K) * (a/n))) K
    using has-integral-spike-finite-eq [of {v} K λx. f (Inf K) f1] f1 by simp
  have f2: f2 x = f (Sup K) if x ∈ K − {u} for x
  proof -
    have x ∈ {u..v}
      using that Kuv greaterThanAtMost-eq-atLeastAtMost-diff by blast
      then have [x * real-of-int n / a] = 1 + int k
        using ⟨n > 0⟩ ⟨a > 0⟩ by (simp add: field-simps u-def v-def a-seg-def
ceiling-eq-iff)
      then show ?thesis
        by (simp add: mult.commute f2-def upper-def a-seg-def SupK v-def)
  qed
  have ((λx. f (Sup K)) has-integral (f (Sup K) * (a/n))) K
    using ⟨n > 0⟩ ⟨a > 0⟩ has-integral-const-real [of f (Sup K) u v]
      by (simp add: Kuv field-simps u-def v-def a-seg-def)

```

```

then show (f2 has-integral (f (Sup K) * (a/n))) K
  using has-integral-spoke-finite-eq [of {u} K λx. f (Sup K) f2] f2 by simp
  have |v - u| < del (ε/a)
    using <n > 0> <a > 0> by (simp add: v-def u-def a-seg-def field-simps
an-less-del)
    then have |f v - f u| < ε/a
      using <ε > 0> <a > 0> <0 ≤ u> <u ≤ a> <0 ≤ v> <v ≤ a>
      by (intro del) auto
    then show |f(Sup K) - f(Inf K)| < ε/a
      using InfK SupK by blast
qed

have int-21-D: ((λx. f2 x - f1 x) has-integral (f(Sup K) - f(Inf K)) * (a/n))
K if K ∈ ?D for K
  using that has-integral-diff [OF int-f2-D int-f1-D] by (simp add: algebra-simps)

have D-ne: ?D ≠ {}
  by (metis <0 < a> <n > 0> card-gt-0-iff card-regular-division)
have f12: ((λx. f2 x - f1 x) has-integral (∑ K ∈ ?D. (f(Sup K) - f(Inf K)) *
(a/n))) {0..a}
  by (intro div int-21-D has-integral-combine-division)
moreover have (∑ K ∈ ?D. (f(Sup K) - f(Inf K)) * (a/n)) < ε
proof -
  have (∑ K ∈ ?D. (f(Sup K) - f(Inf K)) * (a/n)) ≤ (∑ K ∈ ?D. |f(Sup K) -
f(Inf K)| * (a/n))
    using <n > 0> <a > 0>
    by (smt (verit) divide-pos-pos of-nat-0-less-iff sum-mono zero-le-mult-iff)
  also have ... < (∑ K ∈ ?D. ε/n)
    using <n > 0> <a > 0> less
    by (intro sum-strict-mono finite-regular-division D-ne) (simp add: field-simps)
  also have ... = ε
    using <n > 0> <a > 0> by simp
  finally show ?thesis .
qed

ultimately have f2-near-f1: integral {0..a} (λx. f2 x - f1 x) < ε
  by (simp add: integral-unique)

define yidx where yidx ≡ λy. LEAST k. y < f (a-seg (Suc k))
have fa-yidx-le: f (a-seg (yidx y)) ≤ y and yidx-gt: y < f (a-seg (Suc (yidx
y)))
  if y ∈ {0..b} for y
proof -
  obtain x where x: f x = y x ∈ {0..a}
    using Topological-Spaces.IVT' [OF - - - cont-0a] assms
    by (metis <y ∈ {0..b}> atLeastAtMost-iff)
  define k where k ≡ nat ⌊x/a * n⌋
  have x-lims: a-seg k ≤ x x < a-seg (Suc k)
    using <n > 0> <0 < a> floor-divide-lower floor-divide-upper [of a x*n] x
    by (auto simp: k-def a-seg-def field-simps)

```

```

with that  $x$  obtain  $f$ -lims:  $f(a\text{-seg } k) \leq y \leq f(a\text{-seg } (\text{Suc } k))$ 
  using strict-mono-onD [OF sm] by force
then have  $a\text{-seg } (yidx y) \leq a\text{-seg } k$ 
  by (simp add: Least-le ‹strict-mono a-seg› strict-mono-less-eq yidx-def)
then have  $f(a\text{-seg } (yidx y)) \leq f(a\text{-seg } k)$ 
  using strict-mono-onD [OF sm] by simp
then show  $f(a\text{-seg } (yidx y)) \leq y$ 
  using f-lims by linarith
show  $y < f(a\text{-seg } (\text{Suc } (yidx y)))$ 
  by (metis LeastI f-lims(2) yidx-def)
qed

have yidx-equality:  $yidx y = k$  if  $y \in \{0..b\}$   $y \in \{f(a\text{-seg } k)..f(a\text{-seg } (\text{Suc } k))\}$  for  $y k$ 
proof (rule antisym)
  show  $yidx y \leq k$ 
    unfolding yidx-def by (metis atLeastLessThan-iff that(2) Least-le)
  have  $(a\text{-seg } (\text{real } k)) < a\text{-seg } (1 + \text{real } (yidx y))$ 
    using yidx-gt [OF that(1)] that(2) strict-mono-onD [OF sm] order-le-less-trans
  by fastforce
  then have  $\text{real } k < 1 + \text{real } (yidx y)$ 
    by (simp add: ‹strict-mono a-seg› strict-mono-less)
  then show  $k \leq yidx y$ 
    by simp
qed

have yidx b = n
proof -
  have  $a < (1 + \text{real } n) * a / \text{real } n$ 
    using ‹0 < n› ‹0 < a› by (simp add: divide-simps)
  then have  $b < f(a\text{-seg } (1 + \text{real } n))$ 
    using f ‹a ≥ 0› a-seg-def sm strict-mono-onD by fastforce
  then show ?thesis
    using ‹0 ≤ b› by (auto simp: f a-seg-def yidx-equality)
qed

moreover have yidx-less-n:  $yidx y < n$  if  $y < b$  for  $y$ 
  by (metis ‹0 < n› fa-eq-b gr0-conv-Suc less-Suc-eq-le that Least-le yidx-def)
ultimately have yidx-le-n:  $yidx y \leq n$  if  $y \leq b$  for  $y$ 
  by (metis dual-order.order-iff-strict that)

have zero-to-b-eq:  $\{0..b\} = (\bigcup_{k < n} \{f(a\text{-seg } k)..f(a\text{-seg } (\text{Suc } k))\})$  (is ?lhs = ?rhs)
proof
  show ?lhs ⊆ ?rhs
  proof
    fix y assume y:  $y \in \{0..b\}$ 
    have fn:  $f(a\text{-seg } n) = b$ 
      using a-seg-eq-a-iff ‹f a = b› by fastforce
    show  $y \in ?rhs$ 
    proof (cases y=b)

```

```

case True
with fn <n>0 show ?thesis
    by (rule-tac a=n-1 in UN-I) auto
next
    case False
    with y show ?thesis
        apply (simp add: subset-iff Bex-def)
        by (metis atLeastAtMost-iff of-nat-Suc order-le-less yidx-gt fa-yidx-le
yidx-less-n)
    qed
qed
show ?rhs ⊆ ?lhs
    apply clarsimp
    by (smt (verit, best) a-seg-ge-0 a-seg-le-aff-iff(2) nat-less-real-le of-nat-0-le-iff)
qed

define g1 where g1 ≡ λy. if y=b then a else a-seg (Suc (yidx y))
define g2 where g2 ≡ λy. if y=0 then 0 else a-seg (yidx y)
have g1: g1 y ∈ {0..a} if y ∈ {0..b} for y
    using that <a > 0> yidx-less-n [of y] by (auto simp: g1-def a-seg-def di-
vide-simps)
have g2: g2 y ∈ {0..a} if y ∈ {0..b} for y
    using that <a > 0> yidx-le-n [of y] by (simp add: g2-def a-seg-def divide-simps)

have g2-le-g: g2 y ≤ g y if y ∈ {0..b} for y
proof –
    have f (g2 y) ≤ y
    using <f 0 = 0> g2-def that fa-yidx-le by presburger
    then have f (g2 y) ≤ f (g y)
    using that g by (smt (verit, best) atLeastAtMost-iff fim image-iff)
    then show ?thesis
    by (smt (verit, best) atLeastAtMost-iff fim g g2 imageE sm-0a strict-mono-onD
that)
qed
have g-le-g1: g y ≤ g1 y if y ∈ {0..b} for y
proof –
    have y ≤ f (g1 y)
    by (smt (verit, best) <f a = b> g1-def that yidx-gt)
    then have f (g y) ≤ f (g1 y)
    using that g by (smt (verit, best) atLeastAtMost-iff fim image-iff)
    then show ?thesis
    by (smt (verit, ccfv-threshold) atLeastAtMost-iff f-iff(1) g1 that)
qed

define DN where DN ≡ λK. nat ⌈ Inf K * real n / a ⌉
have [simp]: DN {a * real k / n..a * (1 + real k) / n} = k for k
    using <n > 0> <a > 0> by (simp add: DN-def divide-simps)
have DN: bij-betw DN ?D {..<n}
proof (intro bij-betw-imageI)

```

```

show inj-on DN (regular-division 0 a n)
proof
  fix K K'
  assume K ∈ regular-division 0 a n
  with ⟨a > 0⟩ obtain k where k: K = {a * (real k / n) .. a * (1 + real k) / n}
  by (force elim: regular-divisionE)
  assume K' ∈ regular-division 0 a n
  with ⟨a > 0⟩ obtain k' where k': K' = {a * (real k' / n) .. a * (1 + real k') / n}
  by (force elim: regular-divisionE)
  assume DN K = DN K'
  then show K = K' by (simp add: k k')
qed
have ∃ K ∈ regular-division 0 a n. k = nat ⌈Inf K * real n / a⌉ if k < n for k
  using ⟨n > 0⟩ ⟨a > 0⟩ that
  by (force simp: divide-simps intro: regular-division-eqI [OF refl])
  with ⟨a>0⟩ show DN ‘regular-division 0 a n = {..<n}
  by (auto simp: DN-def bij-betw-def image-iff frac-le elim!: regular-divisionE)
qed

have int-f1: (f1 has-integral (∑ k< n. f(a-seg k)) * (a/n)) {0..a}
proof –
  have a-seg (real (DN K)) = Inf K if K ∈ ?D for K
  using that ⟨a>0⟩ by (auto simp: DN-def field-simps a-seg-def elim: regular-divisionE)
  then have (∑ K ∈ ?D. f(Inf K) * (a/n)) = (∑ k< n. (f(a-seg k)) * (a/n))
  by (simp flip: sum.reindex-bij-betw [OF DN])
  moreover have (f1 has-integral (∑ K ∈ ?D. f(Inf K) * (a/n))) {0..a}
  by (intro div int-f1-D has-integral-combine-division)
  ultimately show ?thesis
  by (metis sum-distrib-right)
qed

The claim (f2 has-integral (∑ k< n. f (a-seg (real (Suc k)))) * (a / real n)) {0..a} can similarly be proved

have int-g1-D: (g1 has-integral a-seg (Suc k) * (f (a-seg (Suc k)) − f (a-seg k)))
  {f(a-seg k)..f(a-seg (Suc k))}
and int-g2-D: (g2 has-integral a-seg k * (f (a-seg (Suc k)) − f (a-seg k)))
  {f(a-seg k)..f(a-seg (Suc k))}
  if k < n for k
proof –
  define u where u ≡ f (a-seg k)
  define v where v ≡ f (a-seg (Suc k))
  obtain u < v 0 ≤ u 0 ≤ v
  unfolding u-def v-def assms
  by (smt (verit, best) a-seg-ge-0 a-seg-le-iff f(1) f-iff(1) of-nat-0-le-iff
of-nat-Suc)

```

```

have  $u \leq b$   $v \leq b$ 
  using  $\langle k < n \rangle \langle a \geq 0 \rangle$  by (simp-all add: u-def v-def flip:  $\langle f a = b \rangle$ )
have  $yidx\text{-eq}: yidx x = k$  if  $x \in \{u..v\}$  for  $x$ 
  using  $\langle 0 \leq u \rangle \langle v \leq b \rangle$  that u-def v-def yidx-equality by auto

have  $g1 x = a\text{-seg} (\text{Suc } k)$  if  $x \in \{u..v\}$  for  $x$ 
  using that  $\langle v \leq b \rangle$  by (simp add: g1-def yidx-eq)
  moreover have  $((\lambda x. a\text{-seg} (\text{Suc } k)) \text{ has-integral } (a\text{-seg} (\text{Suc } k) * (v-u)))$ 
  {u..v}
    using has-integral-const-real  $\langle u < v \rangle$ 
    by (metis content-real-if less-eq-real-def mult.commute real-scaleR-def)
  ultimately show  $(g1 \text{ has-integral } (a\text{-seg} (\text{Suc } k) * (v-u)))$  {u..v}
    using has-integral-spike-finite-eq [of {v} {u..v}]  $\lambda x. a\text{-seg} (\text{Suc } k) g1$  by
simp

have  $g2: g2 x = a\text{-seg } k$  if  $x \in \{u..v\}$  for  $x$ 
  using that  $\langle 0 \leq u \rangle$  by (simp add: g2-def yidx-eq)
  moreover have  $((\lambda x. a\text{-seg } k) \text{ has-integral } (a\text{-seg } k * (v-u)))$  {u..v}
    using has-integral-const-real  $\langle u < v \rangle$ 
    by (metis content-real-if less-eq-real-def mult.commute real-scaleR-def)
  ultimately show  $(g2 \text{ has-integral } (a\text{-seg } k * (v-u)))$  {u..v}
    using has-integral-spike-finite-eq [of {u,v} {u..v}]  $\lambda x. a\text{-seg } k g2$  by simp
qed

have int-g1:  $(g1 \text{ has-integral } (\sum k < n. a\text{-seg} (\text{Suc } k) * (f (a\text{-seg} (\text{Suc } k)) - f (a\text{-seg } k))))$  {0..b}
  and int-g2:  $(g2 \text{ has-integral } (\sum k < n. a\text{-seg } k * (f (a\text{-seg} (\text{Suc } k)) - f (a\text{-seg } k))))$  {0..b}
    unfolding zero-to-b-eq using int-g1-D int-g2-D
    by (auto simp: min-def pairwise-def intro!: has-integral-UN negligible-atLeastAtMostI)

have  $(\sum k < n. a\text{-seg} (\text{Suc } k) * (f (a\text{-seg} (\text{Suc } k)) - f (a\text{-seg } k)))$ 
   $= (\sum k < n. (\text{Suc } k) * (f (a\text{-seg} (\text{Suc } k)) - f (a\text{-seg } k))) * (a/n)$ 
  unfolding a-seg-def sum-distrib-right sum-divide-distrib by (simp add: mult-ac)
  also have ...  $= (n * f (a\text{-seg } n) - (\sum k < n. f (a\text{-seg } k))) * a / n$ 
    using weighted-nesting-sum [where  $g = f \circ a\text{-seg}$ ] by simp
  also have ...  $= a * b - (\sum k < n. f (a\text{-seg } k)) * a / n$ 
    using <math>n > 0</math> by (simp add: fa-eq-b field-simps)
  finally have int-g1':  $(g1 \text{ has-integral } a * b - (\sum k < n. f (a\text{-seg } k)) * a / n)$ 
  {0..b}
    using int-g1 by simp

The claim  $(g2 \text{ has-integral } a * b - (\sum k < n. f (a\text{-seg} (\text{real } (\text{Suc } k)))) * a / \text{real } n)$  {0..b} can similarly be proved.

have a-seg-diff:  $a\text{-seg} (\text{Suc } k) - a\text{-seg } k = a/n$  for  $k$ 
  by (simp add: a-seg-def field-split-simps)
have f-a-seg-diff:  $|f (a\text{-seg} (\text{Suc } k)) - f (a\text{-seg } k)| < \varepsilon/a$  if  $k < n$  for  $k$ 
  using that <math>a > 0</math> a-seg-diff an-less-del <math>\varepsilon > 0</math>
  by (intro del) auto

```

```

have (( $\lambda x. g1 x - g2 x$ ) has-integral ( $\sum k < n. (f(a\text{-seg}(Suc k)) - f(a\text{-seg} k)) * (a/n)$ )
*  $\{0..b\}$ 
  using has-integral-diff [OF int-g1 int-g2] a-seg-diff
  apply (simp flip: sum-subtractf left-diff-distrib)
  apply (simp add: field-simps)
  done
moreover have ( $\sum k < n. (f(a\text{-seg}(Suc k)) - f(a\text{-seg} k)) * (a/n) < \varepsilon$ )
proof -
  have ( $\sum k < n. (f(a\text{-seg}(Suc k)) - f(a\text{-seg} k)) * (a/n)$ )
     $\leq (\sum k < n. |f(a\text{-seg}(Suc k)) - f(a\text{-seg} k)| * (a/n))$ 
    by simp
  also have ...  $< (\sum k < n. (\varepsilon/a) * (a/n))$ 
  proof (rule sum-strict-mono)
    fix k assume k  $\in \{.. < n\}$ 
    with  $\langle n > 0 \rangle \langle a > 0 \rangle$  divide-strict-right-mono f-a-seg-diff pos-less-divide-eq
    show  $|f(a\text{-seg}(Suc k)) - f(a\text{-seg} k)| * (a/n) < \varepsilon/a * (a/n)$  by fastforce
  qed (use  $\langle n > 0 \rangle$  in auto)
  also have ...  $= \varepsilon$ 
    using  $\langle n > 0 \rangle \langle a > 0 \rangle$  by simp
  finally show ?thesis .
qed
ultimately have g2-near-g1: integral  $\{0..b\} (\lambda x. g1 x - g2 x) < \varepsilon$ 
  by (simp add: integral-unique)

have ab1: integral  $\{0..a\} f1 + integral \{0..b\} g1 = a*b$ 
  using int-f1 int-g1' by (simp add: integral-unique)

have integral  $\{0..a\} (\lambda x. f x - f1 x) \leq integral \{0..a\} (\lambda x. f2 x - f1 x)$ 
proof (rule integral-le)
  show ( $\lambda x. f x - f1 x$ ) integrable-on  $\{0..a\}$  ( $\lambda x. f2 x - f1 x$ ) integrable-on
 $\{0..a\}$ 
  using Henstock-Kurzweil-Integration.integrable-diff int-f1 intgb-f f12 by
blast+
  qed (auto simp: f2-upper)
  with f2-near-f1 have integral  $\{0..a\} (\lambda x. f x - f1 x) < \varepsilon$ 
    by simp
  moreover have integral  $\{0..a\} f1 \leq integral \{0..a\} f$ 
    by (intro integral-le has-integral-integral intgb-f has-integral-integrable [OF
int-f1])
    (simp add: f1-lower)
  ultimately have f-error:  $|integral \{0..a\} f - integral \{0..a\} f1| < \varepsilon$ 
  using Henstock-Kurzweil-Integration.integral-diff int-f1 intgb-f by fastforce

have integral  $\{0..b\} (\lambda x. g1 x - g x) \leq integral \{0..b\} (\lambda x. g1 x - g2 x)$ 
proof (rule integral-le)
  show ( $\lambda x. g1 x - g x$ ) integrable-on  $\{0..b\}$  ( $\lambda x. g1 x - g2 x$ ) integrable-on
 $\{0..b\}$ 
  using Henstock-Kurzweil-Integration.integrable-diff int-g1 int-g2 intgb-g by

```

```

blast+
qed (auto simp: g2-le-g)
with g2-near-g1 have integral {0..b} ( $\lambda x. g1 x - g x$ ) <  $\varepsilon$ 
  by simp
moreover have integral {0..b}  $g \leq \text{integral } \{0..b\} g1$ 
  by (intro integral-le has-integral-integral intgb-g has-integral-integrable [OF
int-g1])
  (simp add: g-le-g1)
ultimately have g-error: |integral {0..b} g1 - integral {0..b} g| <  $\varepsilon$ 
  using integral-diff int-g1 intgb-g by fastforce
show ?thesis
  using f-error g-error ab1 by linarith
qed
show ?thesis
  using * [of | $a * b - \text{integral } \{0..a\} f - \text{integral } \{0..b\} g| / 2] by fastforce
qed (use assms in force)$ 
```

**corollary** Youngs-strict:

```

fixes f :: real  $\Rightarrow$  real
assumes sm: strict-mono-on {0..} f and cont: continuous-on {0..} f and a>0
b $\geq 0$ 
  and f: f 0 = 0 f a  $\neq$  b and fim: f ' $\{0..\}$  = {0..}
  and g:  $\bigwedge x. 0 \leq x \implies g(f x) = x$ 
shows a*b < integral {0..a} f + integral {0..b} g
proof -
have f-iff [simp]: f x < f y  $\longleftrightarrow$  x < y f x  $\leq$  f y  $\longleftrightarrow$  x  $\leq$  y
  if x  $\geq 0$  y  $\geq 0$  for x y
  using that by (smt (verit, best) atLeast-iff sm strict-mono-onD) +
let ?b' = f a
have ?b'  $\geq 0$ 
  by (smt (verit, best) <0 < a atLeast-iff f sm strict-mono-onD)
then have sm-gx: strict-mono-on {0..} g
  unfolding strict-mono-on-def
  by (smt (verit, best) atLeast-iff f-iff(1) f-inv-into-f fim g inv-into-into)
show ?thesis
proof (cases ?b' < b)
case True
have gt-a: a < g y if y  $\in \{?b' < ..b\}$  for y
proof -
have a = g ?b'
  using <a > 0> g by force
also have ... < g y
  using <0  $\leq ?b'\}$  sm-gx strict-mono-onD that by fastforce
finally show ?thesis .
qed
have continuous-on {0..} g
  by (metis cont f(1) fim g sm strict-mono-continuous-invD)

```

```

then have contg: continuous-on {?b'..b} g
  by (meson Icc-subset-Ici-iff ‹0 ≤ f a› continuous-on-subset)
have mono-on {0..} g
  by (simp add: sm-gx strict-mono-on-imp-mono-on)
then have int-g0b: g integrable-on {0..b}
  by (simp add: integrable-on-mono-on mono-on-subset)
then have int-gb'b: g integrable-on {?b'..b}
  by (simp add: ‹0 ≤ ?b'› integrable-on-subinterval)
have a * (b - ?b') = integral {?b'..b} (λy. a)
  using True by force
also have ... < integral {?b'..b} g
  using contg True gt-a by (intro integral-less-real) auto
finally have *: a * (b - ?b') < integral {?b'..b} g .
have a*b = a * ?b' + a * (b - ?b')
  by (simp add: algebra-simps)
also have ... = integral {0..a} f + integral {0..?b'} g + a * (b - ?b')
  using Youngs-exact ‹a > 0› cont ‹f 0 = 0› g sm by force
also have ... < integral {0..a} f + integral {0..?b'} g + integral {?b'..b} g
  by (simp add: *)
also have ... = integral {0..a} f + integral {0..b} g
  by (smt (verit) Henstock-Kurzweil-Integration.integral-combine True ‹0 ≤ ?b'›
int-g0b)
  finally show ?thesis .
next
case False
with f have b < ?b' by force
obtain a' where f a' = b a' ≥ 0
  using fim ‹b ≥ 0› by force
then have a' < a
  using ‹b < f a› ‹a > 0› by force
have gt-b: b < f x if x ∈ {a' <..a} for x
  using ‹0 ≤ a'› ‹f a' = b› that by fastforce
have int-f0a: f integrable-on {0..a}
  by (simp add: integrable-on-mono-on mono-on-def)
then have int-fa'a: f integrable-on {a'..a}
  by (simp add: ‹0 ≤ a'› integrable-on-subinterval)
have cont-f': continuous-on {a'..a} f
  by (meson Icc-subset-Ici-iff ‹0 ≤ a'› cont continuous-on-subset)
have b * (a - a') = integral {a'..a} (λx. b)
  using ‹a' < a› by simp
also have ... < integral {a'..a} f
  using cont-f' ‹a' < a› gt-b by (intro integral-less-real) auto
finally have *: b * (a - a') < integral {a'..a} f .
have a*b = a' * b + b * (a - a')
  by (simp add: algebra-simps)
also have ... = integral {0..a'} f + integral {0..b} g + b * (a - a')
  by (simp add: Youngs-exact ‹0 ≤ a'› ‹f a' = b› cont f g sm)
also have ... < integral {0..a'} f + integral {0..b} g + integral {a'..a} f
  by (simp add: *)

```

```

also have ... = integral {0..a} f + integral {0..b} g
  by (smt (verit) Henstock-Kurzweil-Integration.integral-combine ‹0 ≤ a'› ‹a'
< a› int-f0a)
  finally show ?thesis .
qed
qed

corollary Youngs-inequality:
fixes f :: real ⇒ real
assumes sm: strict-mono-on {0..} f and cont: continuous-on {0..} f and a≥0
b≥0
  and f: f 0 = 0 and fim: f ` {0..} = {0..}
  and g: ∀x. 0 ≤ x ⇒ g (f x) = x
shows a*b ≤ integral {0..a} f + integral {0..b} g
proof (cases a=0)
  case True
  have g x ≥ 0 if x ≥ 0 for x
    by (metis atLeast-iff fim g imageE that)
  then have 0 ≤ integral {0..b} g
    by (metis Henstock-Kurzweil-Integration.integral-nonneg atLeastAtMost-iff
        not-integrable-integral order-refl)
  then show ?thesis
    by (simp add: True)
next
  case False
  then show ?thesis
    by (smt (verit) assms Youngs-exact Youngs-strict)
qed

end

```

## References

- [1] F. Cunningham and N. Grossman. On Young's inequality. *The American Mathematical Monthly*, 78(7):781–783, 1971.