# Wooley's Discrete Inequality

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#### Abstract

This is a formalisation of the proof of an inequality by Trevor D. Wooley attesting that when  $\lambda > 0$ ,

$$\min_{r \in \mathbb{N}} (r + \lambda/r) \le \sqrt{4\lambda + 1}$$

with equality if and only if  $\lambda = m(m-1)$  for some positive integer m.

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## 1 Wooley's Discrete Inequality

theory Wooley-Elementary-Discrete-Inequality imports HOL-Library.Quadratic-Discriminant HOL-Real-Asymp.Real-Asymp

#### begin

This is a formalisation of the proof of an inequality by Trevor D. Wooley attesting that when  $\lambda > 0$ ,

$$\min_{r \in \mathbb{N}} (r + \lambda/r) \le \sqrt{4\lambda + 1}$$

with equality if and only if  $\lambda = m(m-1)$  for some positive integer m. The source is the note "An Elementary Discrete Inequality" available on Wooley's webpage [1]: https://www.math.purdue.edu/~twooley/publ/20230410discineq.pdf.

#### 1.1 General elementary technical lemmas

```
\mathbf{lemma}\ obtains-nat-in-interval:
 fixes x::real assumes x>0
 obtains c::nat where c \in \{x < ... x+1\}
proof
 show nat|x+1| \in \{x < ... x + 1\}
   using assms by force
qed
lemma obtains-nat-in-interval-greater-leq:
 fixes x::real assumes x>0
 obtains c::nat where c > x and c < x+1
 by (meson assms greaterThanAtMost-iff obtains-nat-in-interval)
lemma obtains-nat-in-interval-half:
 fixes x::real assumes x \ge 1/2
 obtains c::nat where c > x - (1/2) and c \le x+1/2
 using assms obtains-nat-in-interval-greater-leq [of x-1/2]
 by (smt (verit) field-sum-of-halves)
```

# 1.2 Trivial case, where we minimise over all positive real values of r

```
{\bf theorem}\ \ elementary\text{-}ineq\text{-}Wooley\text{-}real\text{:}
 fixes l::real and g::real \Rightarrow real
 assumes l > 0 and \forall r \in R. g r = r + (l/r)
   and R = \{r :: real. \ r > 0\}
 shows (\forall r \in R. g \ r \geq 2* sqrt(l)) \land (\forall r \in R. g \ (sqrt(l)) \leq g \ r)
proof-
  have \forall r \in R. 2 * sqrt(l) + (sqrt(r) - (sqrt(l)/sqrt(r)))^2 = r + (l/r)
   using assms by (simp add: power-divide power2-diff)
 moreover
 have \forall r \in R. 2 * sqrt(l) + (sqrt(r) - (sqrt(l)/sqrt(r)))^2 \ge 2 * (sqrt(l))
   using assms by auto
 ultimately have \forall r \in R. \ r+(l/r) \ge 2*sqrt(l) by simp
 moreover
 have g(sqrt(l)) = 2 *sqrt(l) using assms by (simp \ add: real-div-sqrt)
  ultimately show ?thesis using assms by auto
qed
```

### 1.3 Main result: Inequality for the discrete version

```
theorem elementary-discrete-ineq-Wooley: fixes l::real and g::nat \Rightarrow real assumes l > 0 and R = \{r::nat. r > 0\} and \forall r \in R. g r = r + (l/r) shows (INF \ r \in R. \ g \ r) \leq sqrt(4*l+1)
```

We will first show the inequality for a specific choice of  $r_u \in R$ . Then the assertion of the theorem will be simply shown by transitivity.

**define** x::real where x = sqrt(l+1/4)

```
with assms have x>1/2
   by (smt (verit, best) real-sqrt-divide real-sqrt-four real-sqrt-less-iff real-sqrt-one)
  obtain r-u::nat where r-u > x - 1/2 and r-u \le x + 1/2
    using obtains-nat-in-interval-half \langle x > 1/2 \rangle by (metis less-eq-real-def)
  have r - u \in R using assms \langle 1 / 2 \langle x \rangle \langle x - 1 / 2 \langle real \ r - u \rangle by auto
  have ru-gt: r-u > sqrt(l+1/4)-1/2 using \langle r-u > x - 1/2 \rangle \langle x = sqrt(l+1/4) \rangle
  have ru-le: r-u \le sqrt(l+1/4)+1/2 using \langle r-u \le x+1/2 \rangle \langle x = sqrt(l+1/4) \rangle
by blast
    Proving the following auxiliary statement is the key part of the whole
proof.
  have auxiliary: |r-u - (l/r-u)| < 1
  proof-
    define \delta::real where \delta = r-u - sqrt(l+1/4)
    with assms ru-gt \delta-def ru-le
    have \delta: \delta > -1/2 \delta \leq 1/2
      by auto
    have a: |r-u - l/r-u| = |((sqrt(l+1/4) + \delta)^2 - l)/(sqrt(l+1/4) + \delta)|
      using \delta-def
      \mathbf{by} \ (\mathit{smt} \ (\mathit{verit}, \ \mathit{ccfv}\text{-}\mathit{SIG}) \ \land 1 \ / \ 2 \ < x \ \land \ \alpha - \ 1 \ / \ 2 \ < \mathit{real} \ \mathit{r-u} \ \rangle
          add-divide-distrib nonzero-mult-div-cancel-right power2-eq-square)
   have b:|((sqrt(l+1/4) + \delta)^2 - l)/(sqrt(l+1/4) + \delta)| =
|2*\delta + (((1/4) - \delta^2)/(sqrt(l+1/4) + \delta))|
    proof-
     have |((sqrt(l+1/4) + \delta)^2 - l)/(sqrt(l+1/4) + \delta)| =
 |(1/4 + 2*(sqrt(l+1/4))*\delta + \delta^2)/(sqrt(l+1/4) + \delta)|
     by (smt (verit, best) assms(1) divide-nonneg-nonneg power2-sum real-sqrt-pow2)
     also have ... = |(2*\delta*(sqrt(l+1/4)) + 2*\delta^2 + 1/4 - \delta^2)/(sqrt(l+1/4))
+\delta
       by (smt (verit) power2-sum)
      also have ... = |(2*\delta*(sqrt(l+1/4)+\delta) + 1/4 - \delta^2)/(sqrt(l+1/4) + \delta)|
\delta)
        by (smt (verit, ccfv-SIG) power2-diff power2-sum)
      also have ... = \frac{(2* \delta* (sqrt(l+1/4)+\delta))}{(sqrt(l+1/4)+\delta)}
+((1/4 -\delta^2)/(sqrt(l+1/4) + \delta))
        \mathbf{by} \ (\mathit{metis} \ \mathit{add-diff-eq} \ \mathit{add-divide-distrib})
      also have ... = |2*\delta + ((1/4 - \delta^2)/(sqrt(l+1/4) + \delta))|
        using \langle \delta = real \ r - u - sqrt \ (l + 1 \ / \ 4) \rangle \langle r - u \in R \rangle assms by force
      finally show ?thesis.
    qed
    show ?thesis
```

```
We distinguish the cases \delta > 0 and \delta < 0:
   proof (cases \delta > 0)
     {f case}\ True
     define t::real where t = 1/2 - \delta
     have c: 0 \le 2*\delta + ((1/4 - \delta^2)/(sqrt(l+1/4) + \delta))
     proof-
       have \delta^2 \le 1/4 using \delta \langle \delta > \theta \rangle
         by (metis less-eq-real-def plus-or-minus-sqrt real-sqrt-divide real-sqrt-four
real-sqrt-le-iff real-sqrt-one real-sqrt-power)
       then have 1/4 - \delta^2 \ge 0
         by simp
       then show ?thesis using \langle \delta > 0 \rangle assms by simp
      have d: 2*\delta + ((1/4 - \delta^2)/(sqrt(l+1/4) + \delta)) \le 1 - 2*t + ((t-t^2)/(sqrt(l+1/4) + \delta))
(1-t)
     proof-
       have \delta = 1/2 - t using t-def by simp
       then have 2*\delta + ((1/4 - \delta^2)/(sqrt(l+1/4) + \delta)) =
                 2*(1/2-t)+((1/4-(1/2-t)^2)/(sqrt(l+1/4)+1/2-t)
))
         by simp
     also have \dots = 1 - 2*t + ((1/4 - (1/4 - 2*(1/2)*t + t^2))/(sqrt(l+1/4))
+1/2-t)
         by (simp add: power2-diff power-divide)
       also have ... = 1 - 2*t + ((t-t^2)/(sqrt(l+1/4) + 1/2 - t)) by simp
       also have ... \leq 1 - 2*t + ((t-t^2)/(1-t))
       proof-
         have sqrt(l+1/4) + 1/2 \ge 1
           using \langle 1/2 \langle x \rangle x-def by linarith
         then have *: sqrt(l+1/4) + 1/2 - t \ge 1 - t by simp
         have 1-t \neq 0 using \langle t = 1/2 - \delta \rangle \langle \delta > 0 \rangle by linarith
         have sqrt(l+1/4) + 1/2 - t \neq 0
           using \delta-def \langle t = 1/2 - \delta \rangle \langle \delta > 0 \rangle assms(1) by force
         then have (1/(sqrt(l+1/4) + 1/2 - t)) \le (1/(1-t))
               using * \langle 1-t \neq 0 \rangle by (smt (verit) True \langle \delta = 1/2 - t \rangle frac-le
le-divide-eq-1-pos)
         have t-t^2 \ge 0 using \langle \delta = 1/2 - t \rangle \langle \delta > 0 \rangle
          by (smt\ (verit,\ best)\ \langle\delta\leq 1\ /\ 2\rangle\ field-sum-of-halves le-add-same-cancel1
nat-1-add-1
               power-decreasing-iff
          power-one-right real-sqrt-pow2-iff real-sqrt-zero zero-less-one-class.zero-le-one)
         then have ((t-t^2)/(sqrt(t+1/4)+1/2-t)) \le ((t-t^2)/(1-t))
           by (smt\ (verit) * True \langle t = 1/2 - \delta \rangle \ frac-le\ le-divide-eq-1-pos)
         then show ?thesis by force
       qed
       finally show ?thesis.
     qed
```

```
have e: 1 - 2*t + ((t-t^2)/(1-t)) \le 1
     proof-
      have 1 - 2*t + ((t-t^2)/(1-t)) = 1 - 2*t + ((1-t)*t/(1-t)) by algebra
      also have \dots = 1 - t
        using c d bv fastforce
      finally show ?thesis
        using \delta t-def by linarith
     qed
     show ?thesis using a b c d e by linarith
   next
     {f case} False
     define t::real where t = 1/2 + \delta
     then have \delta = t - 1/2 by simp
     have \delta < \theta using False by auto
     have -(2*\delta + ((1/4) - \delta^2)/(sqrt(l+1/4) + \delta))) =
-(2*(t-1/2)+(((1/4)-(t-1/2)^2)/(sqrt(l+1/4)+t-1/2)))
      using \langle \delta = t - 1/2 \rangle by auto
     also have ... = -(2*t-1+((t-t^2)/(sqrt(t+1/4)+t-1/2)))
      by (simp add: power2-diff power-divide)
     finally have ***: -(2*\delta + ((1/4) - \delta^2)/(sqrt(l+1/4) + \delta))) =
-(2*t-1+((t-t^2)/(sqrt(l+1/4)+t-1/2))).
    have c:-(2*\delta+((1/4)-\delta^2)/(sqrt(l+1/4)+\delta))) \le 1-2*t-((t-t^2)/(sqrt(l+1/4)+\delta))
(sqrt(l+1/4)))
     proof-
      have c1: sqrt(l+1/4) + t - 1/2 \le sqrt(l+1/4)
        using \langle \delta = t - 1 / 2 \rangle \langle \delta \leq 0 \rangle by simp
      have (sqrt(l+1/4) + t - 1/2) \neq 0 sqrt(l+1/4) \neq 0
        using assms \delta-def \langle \delta = t - 1/2 \rangle \langle r-u \in R \rangle by auto
      have c2: (t-t^2)/(sqrt(l+1/4) + t - 1/2) \ge (t-t^2)/(sqrt(l+1/4))
        using c1 assms
        by (smt (verit, best) \delta-def ru-gt \langle t = 1/2 + \delta \rangle
           field-sum-of-halves frac-le le-add-same-cancel1 nat-1-add-1 of-nat-0-le-iff
            power-decreasing-iff power-one-right zero-less-one-class.zero-le-one)
      have c3: -(t-t^2)/(sqrt(l+1/4) + t - 1/2) \le -(t-t^2)/(sqrt(l+1/4))
        using c2 by linarith
      show ?thesis using *** c3 by linarith
     qed
```

```
have d: 1 - 2*t - ((t-t^2)/(sqrt(l+1/4))) \le 1
     proof-
       have *: t > 0 using \langle \delta \rangle - 1/2 \rangle \langle t = 1/2 + \delta \rangle by simp
       have **: t \le 1 using \langle \delta \le 0 \rangle \langle t = 1/2 + \delta \rangle by simp
       show ?thesis using * **
             by (smt (verit) assms(1) divide-nonneg-nonneg mult-le-cancel-right2
power2-eq-square real-sqrt-ge-0-iff)
    have e: -(2*\delta + ((1/4) - \delta^2)/(sqrt(l+1/4) + \delta))) \ge 1 - 2*t - ((t-t^2)/t)
       have -(2*\delta + ((1/4) - \delta^2)/(sqrt(l+1/4) + \delta)))
= -(2*t-1+((t-t^2)/(sqrt(t+1/4)+t-1/2)))
         using *** by simp
       have ((t-t^2)/(sqrt(l+1/4)+t-1/2)) \le (t-t^2)/t
       proof-
         have \dagger: (sqrt(l+1/4) + t - 1/2) \ge t using assms
          by (smt (verit, best) one-power2 power-divide real-sqrt-four real-sqrt-pow2
sqrt-le-D)
         moreover have t > 0 using \langle \delta \rangle - 1/2 \rangle \langle t = 1/2 + \delta \rangle by simp
         ultimately have (sqrt(l+1/4) + t - 1/2) > 0
           by auto
         show ?thesis using \dagger \langle (sqrt(l+1/4) + t - 1/2) \rangle > 0 \rangle
             \langle \theta < t \rangle
           by (smt\ (verit,\ best)\ \langle \delta \leq 0 \rangle\ \langle t = 1/2 + \delta \rangle
          frac-le le-add-same-cancel1 le-divide-eq-1-pos nat-1-add-1 power-decreasing-iff
               power-one-right zero-less-one-class.zero-le-one)
       qed
       with *** show ?thesis by linarith
     qed
     have f: 1 - 2*t - ((t-t^2)/t) \ge -1/2
       have t > 0 using \langle \delta \rangle - 1/2 \rangle \langle t = 1/2 + \delta \rangle by simp
       then have 1 - 2 * t - ((t-t^2)/t) = 1 - 2 * t - (1 - t)
         by (metis divide-diff-eq-iff less-irrefl one-eq-divide-iff power2-eq-square)
       also have \dots = -t by auto
       finally show ?thesis
         using \langle \delta \leq \theta \rangle \langle t = 1/2 + \delta \rangle by linarith
     show ?thesis using a b c d e f by linarith
   qed
  qed
```

The next step is to show that by the statement named "auxiliary" shown above, we can directly show the desired inequality for the specific  $r_u \in R$ :

```
have (r\cdot u-l/r\cdot u)^2 \le 1 using auxiliary abs-square-le-1 by blast then have (r\cdot u^2 - 2*r\cdot u*(l/r\cdot u) + l^2/r\cdot u^2 \le 1) using power2-diff power-divide assms by (smt\ (verit)\ mult-2\ of\text{-nat-add}\ of\text{-nat-eq-of-nat-power-cancel-iff}) then have r\cdot u^2 - 2*l + l^2/r\cdot u^2 \le 1 using assms (r\cdot u \in R) by force then have r\cdot u^2 + 2*l + l^2/r\cdot u^2 \le (4*l+1) by argo then have r\cdot u^2 + 2*r\cdot u*(l/r\cdot u) + l^2/r\cdot u^2 \le (4*l+1) using assms by simp then have (r\cdot u+(l/r\cdot u))^2 \le (4*l+1) by (smt\ (verit,\ best)\ mult-2\ of\text{-nat-add}\ of\text{-nat-power-eq-of-nat-cancel-iff}\ power2\text{-sum} power\text{-divide}) then have (r\cdot u+(l/r\cdot u)) \le sqrt(4*l+1) using real\text{-le-rsqrt} by blast moreover
```

The following shows that it is enough that we showed the inequality for the specific  $r_u \in R$ , as the statement of the theorem will then simply hold by transitivity.

```
have (INF\ r\in R.\ g\ r)\leq g\ r-u \operatorname{proof}- have bdd\text{-}below\ (g\ 'R) unfolding bdd\text{-}below\text{-}def using assms\ image\text{-}iff by (metis\ add\text{-}increasing\ assms(1)\ divide\text{-}nonneg\text{-}nonneg\ image\text{-}iff\ less-eq\text{-}real\text{-}def} of\text{-}nat\text{-}0\text{-}le\text{-}iff) show\ ?thesis by (simp\ add:\ \langle bdd\text{-}below\ (g\ 'R)\rangle\ \langle r\text{-}u\in R\rangle\ cINF\text{-}lower) ed ultimately show ?thesis\ using\ assms\ \langle r\text{-}u\in R\rangle\ by\ force ed
```

#### 1.4 Special case: Equality for the discrete version

We will now show a special case of the main result where equality holds instead of inequality.

We will need to make use of the following technical lemma, which will be used so as to guarantee that there exists a  $p \in R$  for which the INF of g(r) equals to g(p). To this end, we will show that here the infimum INF can be identified with the minimum Min by restricting to a finite set. As the operator Min in Isabelle is used for finite sets and R is infinite, we used INF in the original formulation, however here Min and INF can be identified.

The following technical lemma is by Larry Paulson:

```
lemma restrict-to-min: fixes l::real and g::nat \Rightarrow real assumes l>0 and R-def: R=\{r::nat.\ r>0\} and g-def: \forall\ r.\ g\ r=r+(l/r) obtains F where finite\ F\ F\subseteq R\ (INF\ r\in R.\ g\ r)=Min\ (g\ 'F)\ F\neq \{\}
```

```
proof -
 have ge\theta: g r \ge \theta for r
   using \langle l > 0 \rangle R-def g-def by (auto simp: g-def)
  then have bdd: bdd-below (g 'R)
   by (auto simp add: q-def R-def bdd-below-def)
 have \forall_F n in sequentially. g \ 1 < g \ n
   by (simp add: g-def) real-asymp
  then obtain N where N > 0 and N: \bigwedge r. r \ge N \implies g \mid 1 < g \mid r
   by (metis Suc-leD eventually-sequentially less-Suc-eq-0-disj)
  define F where F = R \cap \{..N\}
 have F: finite F F \subseteq R
   by (auto simp add: F-def)
 have F \neq \{\}
   using F-def R-def \langle 0 < N \rangle by blast
 have (INF \ r \in R. \ g \ r) = (INF \ r \in F. \ g \ r)
 proof (intro order.antisym cInf-mono bdd)
   show bdd-below (q 'F)
     by (meson ge0 bdd-belowI2)
  next
   \mathbf{fix} \ b
   assume b \in g ' R
   then show \exists a \in g 'F. a \leq b
     unfolding image-iff F-def R-def Bex-def
    by (metis N linorder-not-less IntI atMost-iff mem-Collect-eq nle-le zero-less-one)
  qed (use \langle F \subseteq R \rangle \langle \theta < N \rangle in \langle auto \ simp: R-def \ F-def \rangle)
 also have \dots = Min(g'F)
   using \langle F \neq \{\} \rangle by (simp add: \langle finite F \rangle cInf-eq-Min)
 finally have (INF \ r \in R. \ g \ r) = Min \ (g \ 'F).
  with F show thesis
   using that \langle F \neq \{\} \rangle by blast
qed
    We will make use of the following calculation, which is convenient to
formulate separately as a lemma.
\mathbf{lemma}\ elementary\text{-}discrete\text{-}ineq\text{-}Wooley\text{-}quadratic\text{-}eq\text{-}sol\text{:}
  fixes l::real and g::nat \Rightarrow real
 assumes l>0 and \forall r. g r = r + (l/r) and g r = sqrt(4*l+1)
  shows (r = 1/2 + (1/2)* \ sqrt(\ 4*l + 1)) \lor (r = -1/2 + (1/2)* \ sqrt(4*l + 1))
+1))
proof-
 have eq\theta: r^2 - r*(sqrt(4*l+1)) + l = 0
 proof-
   have r*(r + l/r) = r*(sqrt(4*l+1)) using assms by simp
   then have r^2 + r*(l/r) = r*(sqrt(4*l+1))
     by (simp add: distrib-left power2-eq-square)
   then show ?thesis
    by (smt (verit, ccfv-threshold) assms divide-eq-eq mult.commute real-sqrt-gt-1-iff)
  qed
```

Solving the above quadratic equation gives the following two roots:

```
have roots: (r = 1/2 + (1/2) * sqrt(4 * l + 1)) \lor (r = -1/2 + (1/2) * sqrt(4 * l + 1))
 proof-
   define a::real where a = 1
   define b::real where b = - sqrt(4*l+1)
   define c::real where c = l
   have a*r^2 + b*r + c = 0 using eq0 by (simp add: mult.commute a-def b-def
c-def)
    then have A: (r = (-b + sqrt(discrim a b c))/2*a) \lor (r = (-b - sqrt(discrim a b c))/2*a)
discrim \ a \ b \ c))/\ 2*a)
     using discriminant-iff[of\ a\ r]\ a-def\ by\ simp
   have discrim a b c = b^2 - 4*a*c
     using discrim-def by simp
    then have B: (r = (-b + sqrt(b^2 - 4*a*c))/2*a) \lor (r = (-b - sqrt(b^2 - 4*a*c))/2*a)
-4*a*c))/2*a)
     using A by auto
  then have C: (r = (-b + sqrt(b^2 - 4*c))/2) \vee (r = (-b - sqrt(b^2 - 4*c))/2)
2)
     using a-def by simp
   have b^2 - 4 * c = 1 using b-def c-def assms(1) by auto
   then have (r = (-b + 1)/2) \vee (r = (-b - 1)/2)
     using C by auto
   then show ?thesis using b-def by auto
 qed
 show ?thesis using roots by simp
\mathbf{qed}
    The special case with equality involves a double implication (iff), and we
start by showing one direction.
\textbf{theorem} \ \ elementary\text{-} \textit{discrete-ineq-Wooley-special-case-1}:
 fixes l::real and g::nat \Rightarrow real assumes l > \theta and R = \{r::nat. \ r > \theta\} and \forall r. g
r = r + (l/r)
   and (INF \ r \in R. \ g \ r) = sqrt(4*l+1)
 shows \exists m::nat. l = m*(m-1)
proof-
 have \exists p \in R. (INF \ r \in R. \ g \ r) = g \ p
 proof-
   obtain F where *:\langle (INF \ r \in R. \ g \ r) = Min \ (g \ `F) \rangle and \langle finite \ F \rangle and \langle F \rangle
\subseteq R \land \langle F \neq \{\} \rangle
     using assms restrict-to-min by metis
   then obtain p::nat where Min(g'F) = gp p \in R
     by (smt (verit) Min-in finite-imageI image-iff image-is-empty subsetD)
   with * show ?thesis by metis
  qed
  with assms
  obtain r-u::nat where g r-u = sqrt(4*l+1) and r-u \in R
   by metis
  then have ru: (r-u + (l/r-u)) = sqrt( 4*l+1 )
   using assms by auto
```

```
have (r-u = 1/2 + (1/2)* sqrt(4*l+1)) \Longrightarrow (l = r-u^2 - r-u)
 proof-
   assume r - u = 1/2 + (1/2) * (sqrt(4 * l + 1))
   then have 2*r-u = 1 + sqrt(4*l + 1) by simp
   then have (2* real(r-u) - 1)^2 = (4*l + 1) using assms by auto
   then have (2*real(r-u))^2 - 2*(2*real(r-u)) + 1 = (4*l+1)
    by (simp add: power2-diff)
   then have 4*real(r-u)^2-4*(r-u)=4*l by fastforce
   then show (l = r - u^2 - r - u)
    by (simp add: of-nat-diff power2-eq-square)
 moreover
 have (r-u = -1/2 + (1/2)* sqrt(4*l+1)) \Longrightarrow (l = r-u^2 + r-u)
 proof-
   assume r - u = -1/2 + (1/2) * sqrt(4 * l + 1)
   then have 2 * r-u +1 = sqrt(4*l+1) by simp
   then have (2*r-u+1)^2 = (4*l+1) using assms by auto
   then have 4*(r-u)^2 + 4*r-u + 1 = 4*l+1
    by (simp add: power2-eq-square)
   then show (l = r - u^2 + r - u)
    by (simp add: of-nat-diff power2-eq-square)
 qed
 moreover
 have (r-u = 1/2 + (1/2)* sqrt(4*l+1)) \lor (r-u = -1/2 + (1/2)* sqrt(4*l+1))
   using assms ru elementary-discrete-ineq-Wooley-quadratic-eq-sol
    assms by auto
 ultimately have (l = r - u^2 + r - u) \lor (l = r - u^2 - r - u)
   \mathbf{by} blast
 then show ?thesis
  by (metis add-implies-diff distrib-left mult.commute mult.right-neutral power2-eq-square
right-diff-distrib')
    (Interestingly, the above use of metis finished the proof in a simple step
guaranteeing the existence of a witness with the desired property).
qed
   Now we show the other direction.
\textbf{theorem}\ elementary-discrete-ineq-Wooley-special-case-2}:
 fixes l::real and g::nat \Rightarrow real
 assumes l>0 and R=\{r::nat.\ r>0\} and \forall\ r.\ g\ r=r+\ (l/r) and \exists\ m::nat.\ l
=m*(m-1)
 shows (INF r \in R. q(r) = sqrt(4*l+1)
proof-
 obtain r-u::nat where (l = r-u^2 + r-u) using assms
   by (metis add.commute add-cancel-left-right mult-eq-if power2-eq-square)
```

```
then have sqrt(4*l+1) = sqrt(4*r-u^2 + 4*r-u + 1) by simp moreover have 4*r-u^2 + 4*r-u + 1 = (2*r-u + 1)^2 by (simp\ add:\ Groups.mult-ac(2)\ distrib-left\ power2-eq-square) ultimately have 4:\ sqrt(4*l+1) = sqrt((2*r-u + 1)^2) by metis then have ru:\ r-u = -1/2 + 1/2*sqrt(4*l+1) by (simp\ add:\ add-divide-distrib)
```

To prove the conclusion of the theorem, we will follow a proof by contradiction.

```
show ?thesis
 proof (rule ccontr)
   assume Inf (g \cdot R) \neq sqrt (4 * l + 1)
   then have inf: (INF \ r \in R. \ g \ r) < sqrt(4*l+1)
     using assms less-eq-real-def elementary-discrete-ineq-Wooley by blast
   have \exists p \in R. (INF \ r \in R. \ g \ r) = g \ p
   proof-
      obtain F where *:\langle (INF \ r \in R. \ g \ r) = Min \ (g \ `F) \rangle and \langle finite \ F \rangle \ \langle F \subseteq F \rangle
R \land \langle F \neq \{\} \rangle
       using assms restrict-to-min by metis
     then obtain p::nat where Min(g'F) = g p p \in R
       by (meson Min-in finite-imageI imageE image-is-empty subsetD)
     with * show ?thesis by metis
   qed
   obtain p::nat where p \in R and (INF \ r \in R. \ g \ r) = g \ p using assms
       \forall \exists p \in R. (INF \ r \in R. \ g \ r) = g \ p \land \mathbf{by} \ blast
   then have (p+l/p < sqrt(4*l+1))
     using inf \ assms(3) by auto
   have p*(p+l/p) < p*(sqrt(4*l+1))
     using \langle p \in R \rangle \langle (p+l/p < sqrt(4*l+1)) \rangle assms by simp
   then have p^2 - p*(sqrt(4*l+1)) + l < 0
    by (smt\ (verit)\ \langle p\in R\rangle\ assms(2)\ distrib-left\ mem-Collect-eq\ nonzero-mult-div-cancel-left
         of-nat-0-less-iff of-nat-mult power2-eq-square times-divide-eq-right)
```

We now need to find the possible values of this hypothetical  $p \in R$ , i.e. the roots of the above quadratic inequality. (These will be in-between the roots of the corresponding quadratic equation which were given in lemma  $\llbracket \theta < ?l; \forall r. ?g \ r = real \ r + ?l \ / \ real \ r; ?g \ ?r = sqrt \ (4 * ?l + 1) \rrbracket \Longrightarrow real ?r = 1 \ / \ 2 + 1 \ / \ 2 * sqrt \ (4 * ?l + 1) \lor real ?r = -1 \ / \ 2 + 1 \ / \ 2 * sqrt \ (4 * ?l + 1)$ ). Here we show that the roots of the quadratic inequality lie in the following interval via a direct calculation:

```
have p: (p < (sqrt(4*l+1) + 1) / 2) \land (p > (sqrt(4*l+1) - 1) / 2) proof — have p^2 - p*(sqrt(4*l+1)) + l + 1/4 < 1/4 using (p^2 - p*(sqrt(4*l+1)) + l < 0) by simp moreover
```

```
have -(2*(p*sqrt(4*l+1))/2) + (4*l+1)/4 = -p*(sqrt(4*l+1))+l
+1/4
                     by force
              ultimately have ***: p^2 - (2*(p* sqrt(4*l+1))/2) + (4*l+1)/4 < 1/4
                     by linarith
                 have ***: (p - (sqrt(4*l+1))/2)^2 = p^2 - 2 * p * (sqrt(4*l+1))/2 + (sqrt(4*l+1))/
(sqrt(4*l+1))/2)^2
                      by (simp add: power2-diff)
               then have p^2 - 2 * p * (sqrt(4*l+1))/2 + ((sqrt(4*l+1))/2)^2 = p^2 - 2
* p* (sqrt(4*l+1))/2 + (4*l+1)/4
                     by (smt (verit) assms(1) power-divide real-sqrt-four real-sqrt-pow2)
                then have (p - (sqrt(4*l + 1))/2)^2 < 1/4 using *** **** by linarith
                then have |(p - (sqrt(4*l + 1))/2)| < 1/2
                              by (metis real-sqrt-abs real-sqrt-divide real-sqrt-four real-sqrt-less-mono
real-sqrt-one)
                   then have ((p - (sqrt(4*l + 1))/2)) < 1/2 ((p - (sqrt(4*l + 1))/2)) >
-1/2 by linarith+
               then show ?thesis
                     by force
          ged
```

So p lies in an interval of length strictly less than 1 between two positive integers, but this means that p cannot be a positive integer, which yields the desired contradiction, thus completing the proof:

```
obtain A::nat where A: real\ A = -1/2 + (1/2)* \ sqrt(4*l+1) using ru by blast then show False using 4 p by fastforce qed qed
```

Finally, for convenience and completeness, we state the special case where equality holds formulated with the double implication and moreover including the values for which the INF (i.e. minimum here as we have seen) is attained as previously calculated.

```
theorem elementary-discrete-ineq-Wooley-special-case-iff: fixes l::real and g::nat \Rightarrow real assumes l>0 and R=\{r::nat.\ r>0\} and \forall\ r.\ g\ r=r+(l/r) shows ((INF\ r\in R.\ g\ r)=sqrt(4*l+1))\longleftrightarrow (\exists\ m::nat.\ l=m*(m-1)) and g\ p=sqrt(4*l+1)\longrightarrow (p=1/2+(1/2)*\ sqrt(\ 4*l+1))\lor (p=-1/2+(1/2)*\ sqrt(\ 4*l+1)) using assms elementary-discrete-ineq-Wooley-special-case-1 elementary-discrete-ineq-Wooley-special-case-2 apply blast using assms(1)\ assms(3)\ elementary-discrete-ineq-Wooley-quadratic-eq-sol\ restrict-to-min by auto
```

 $\mathbf{end}$ 

## References

[1] T. D. Wooley. An elementary discrete inequality. https://www.math.purdue.edu/~twooley/publ/20230410 discineq.pdf.