

Fundamentals of Unconstrained Optimization

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Abstract

As formal methods gain traction in machine learning and numerical analysis, the community needs computer-checked proofs of core optimization results. Existing Isabelle libraries still lack a foundational framework for unconstrained optimization. We close this gap with a comprehensive Isabelle/HOL development that formalizes:

- (1) minimizers, strict and isolated local minimizers;
- (2) first- and second-order optimality conditions for scalar functions $f : \mathbb{R} \rightarrow \mathbb{R}$;
- (3) first-order optimality conditions for vector functions $g : \mathbb{R}^n \rightarrow \mathbb{R}$; and
- (4) a worked example showing that the continuous function

$$h(x) = \begin{cases} x^4(\cos(1/x) + 2), & x \neq 0, \\ 0, & x = 0 \end{cases}$$

has a *strict* but *non-isolated* local minimizer at $x = 0$.

The new session `Unconstrained_Optimization` provides sound, reusable foundations for future proof-checking tools and mechanized research in optimization, analysis, and algorithmic correctness.

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1 Auxiliary Facts

theory *Auxiliary-Facts*

imports

Sigmoid-Universal-Approximation.Limits-Higher-Order-Derivatives

begin

1.1 Differentiation Lemmas

lemma *has-derivative-imp*:

fixes $f :: \text{real} \Rightarrow \text{real}$

assumes $(f \text{ has-derivative } f') \text{ (at } x)$

shows $f \text{ differentiable (at } x) \wedge \text{deriv } f \ x = f' \ 1$

proof *safe*

show $f \text{ differentiable at } x$

by $(\text{meson } \text{assms } \text{differentiableI})$

then show $\text{deriv } f \ x = f' \ 1$

by $(\text{metis } \text{DERIV-deriv-iff-real-differentiable } \text{assms } \text{has-derivative-unique}$
 $\text{has-field-derivative-imp-has-derivative } \text{mult.comm-neutral})$

qed

lemma *DERIV-inverse-func*:

assumes $x \neq 0$

shows $\text{DERIV } (\lambda w. 1 / w) \ x \ := \ -1 / x^2$

proof –

have $\text{inverse} = (/) \ (1::'a)$

using *inverse-eq-divide* by *auto*

then show *?thesis*

by $(\text{metis } (\text{no-types}) \ \text{DERIV-inverse } \text{assms } \text{divide-minus-left } \text{numeral-2-eq-2}$
 $\text{power-one-over})$

qed

lemma *power-rule*:

fixes $z :: \text{real}$ **and** $n :: \text{nat}$
shows $\text{deriv } (\lambda x. x \wedge^n) z = (\text{if } n = 0 \text{ then } 0 \text{ else } \text{real } n * z \wedge^{(n-1)})$
by(*subst deriv-pow, simp-all*)

1.1.1 Transfer Lemmas

lemma *has-derivative-transfer-on-ball*:
fixes $f g :: \text{real} \Rightarrow \text{real}$
assumes $\text{eps-gt0}: 0 < \varepsilon$
assumes $\text{eq-on-ball}: \forall y. y \in \text{ball } x \ \varepsilon \longrightarrow f y = g y$
assumes $f\text{-has-deriv}: (f \text{ has-derivative } D) \text{ (at } x)$
shows $(g \text{ has-derivative } D) \text{ (at } x)$
using *centre-in-ball eps-gt0 eq-on-ball f-has-deriv*
has-derivative-transform-within-open **by** *blast*

corollary *field-differentiable-transfer-on-ball*:
fixes $f g :: \text{real} \Rightarrow \text{real}$
assumes $0 < \varepsilon$
assumes $\text{eq-on-ball}: \forall y. y \in \text{ball } x \ \varepsilon \longrightarrow f y = g y$
assumes $f\text{-diff}: f \text{ field-differentiable at } x$
shows $g \text{ field-differentiable at } x$
by (*metis UNIV-I assms dist-commute-lessI*
field-differentiable-transform-within mem-ball)

1.2 Trigonometric Contraction

lemma *cos-contractive*:
fixes $x y :: \text{real}$
shows $|\cos x - \cos y| \leq |x - y|$
proof –
have $|\cos x - \cos y| = |-2 * \sin ((x + y) / 2) * \sin ((x - y) / 2)|$
by (*smt (verit) cos-diff-cos mult-minus-left*)
also have $\dots \leq |\sin ((x + y) / 2)| * (2 * |\sin ((x - y) / 2)|)$
by (*subst abs-mult, force*)
also have $\dots \leq 2 * |\sin ((x - y) / 2)|$
proof –
have $|\sin ((x + y) / 2)| \leq 1$
using *abs-sin-le-one* **by** *blast*
then have $|\sin ((x + y) / 2)| * (2 * |\sin ((x - y) / 2)|) \leq 1 * (2 * |\sin ((x - y) / 2)|)$
by(*rule mult-right-mono, simp*)
then show *?thesis*
by *linarith*
qed
also have $\dots \leq 2 * |(x - y) / 2|$
using *abs-sin-le-one* **by** (*smt (verit, del-insts) abs-sin-x-le-abs-x*)
also have $\dots = |x - y|$
by *simp*
finally show *?thesis*.
qed

lemma *sin-contractive*:
fixes $x\ y :: \text{real}$
shows $|\sin x - \sin y| \leq |x - y|$
proof –
have $|\sin x - \sin y| = |2 * \cos ((x + y) / 2) * \sin ((x - y) / 2)|$
by (*metis (no-types) mult.assoc mult.commute sin-diff-sin*)
also have $\dots \leq |\cos ((x + y) / 2)| * (2 * |\sin ((x - y) / 2)|)$
by (*subst abs-mult, force*)
also have $\dots \leq 2 * |\sin ((x - y) / 2)|$
proof –
have $|\cos ((x + y) / 2)| \leq 1$
using *abs-cos-le-one* **by** *blast*
then have $|\cos ((x + y) / 2)| * (2 * |\sin ((x - y) / 2)|) \leq 1 * (2 * |\sin ((x - y) / 2)|)$
by (*rule mult-right-mono, simp*)
then show *?thesis*
by *linarith*
qed
also have $\dots \leq 2 * |(x - y) / 2|$
using *abs-sin-le-one* **by** (*smt (verit, del-insts) abs-sin-x-le-abs-x*)
also have $\dots = |x - y|$
by *simp*
finally show *?thesis*.
qed

1.3 Algebraic Factorizations

lemma *biquadrate-diff-biquadrate-factored*:
fixes $x\ y :: \text{real}$
shows $y^4 - x^4 = (y - x) * (y^3 + y^2 * x + y * x^2 + x^3)$
proof –
have $y^4 - x^4 = (y^2 - x^2) * (y^2 + x^2)$
by (*metis mult.commute numeral-Bit0 power-add square-diff-square-factored*)
also have $\dots = (y - x) * (y + x) * (y^2 + x^2)$
by (*simp add: power2-eq-square square-diff-square-factored*)
also have $\dots = (y - x) * (y^3 + y^2 * x + y * x^2 + x^3)$
by (*simp add: distrib-left mult.commute power2-eq-square power3-eq-cube*)
finally show *?thesis*.
qed

1.4 Specific Trigonometric Values

lemma *sin-5pi-div-4*: $\sin (5 * \pi / 4) = - (\text{sqrt } 2 / 2)$
proof –
have $5 * \pi / 4 = \pi + \pi / 4$
by *simp*
moreover have $\sin (\pi + x) = - \sin x$ **for** x
by (*simp add: sin-add*)
ultimately show *?thesis*

using *sin-45* by *presburger*
 qed

lemma *cos-5pi-div-4*: $\cos (5 * \pi / 4) = - (\text{sqrt } 2 / 2)$

proof –
 have $5 * \pi / 4 = \pi + \pi / 4$
 by *simp*
 moreover have $\cos (\pi + x) = - \cos x$ for x
 by (*simp add: cos-add*)
 moreover have $\cos (\pi / 4) = \text{sqrt } 2 / 2$
 by (*simp add: real-div-sqrt cos-45*)
 ultimately show *?thesis*
 by *presburger*
 qed

1.5 Local Sign Preservation of Continuous Functions

1.5.1 Local Positivity

lemma *cont-at-pos-imp-loc-pos*:

fixes $g :: \text{real} \Rightarrow \text{real}$ and $x :: \text{real}$
 assumes *continuous (at x) g* and $g x > 0$
 shows $\exists \delta > 0. \forall y. |y - x| < \delta \longrightarrow g y > 0$

proof –
 from *assms* obtain δ where $\delta\text{-pos}: \delta > 0$
 and $\forall y. |y - x| < \delta \longrightarrow |g y - g x| < (g x) / 2$
 using *continuous-at-eps-delta half-gt-zero* by *blast*
 then have $\forall y. |y - x| < \delta \longrightarrow g y > 0$
 by (*smt (verit, best) field-sum-of-halves*)
 then show *?thesis*
 using $\delta\text{-pos}$ by *blast*
 qed

lemma *cont-at-pos-imp-loc-pos'*:

fixes $g :: \text{real} \Rightarrow \text{real}$ and $x :: \text{real}$
 assumes *continuous (at x) g* and $g x > 0$
 shows $\exists \Delta > 0. \forall \delta. 0 < \delta \wedge \delta \leq \Delta \longrightarrow (\forall y. |y - x| < \delta \longrightarrow g y > 0)$

proof –
 from *assms* obtain δ where $\delta\text{-pos}: \delta > 0$ and $H: \forall y. |y - x| < \delta \longrightarrow g y > 0$
 using *cont-at-pos-imp-loc-pos* by *blast*
 have $\forall \delta' \leq \delta. \forall y. |y - x| < \delta' \longrightarrow g y > 0$
 proof *clarify*
 fix $\delta' y :: \text{real}$
 assume $\delta' \leq \delta$ and $|y - x| < \delta'$
 thus $g y > 0$ by (*auto simp: H*)
 qed
 then show *?thesis*
 using $\delta\text{-pos}$ by *blast*
 qed

1.5.2 Local Negativity

lemma *cont-at-neg-imp-loc-neg*:

fixes $g :: \text{real} \Rightarrow \text{real}$ **and** $x :: \text{real}$

assumes *continuous (at x) g* **and** $g\ x < 0$

shows $\exists \delta > 0. \forall y. |y - x| < \delta \longrightarrow g\ y < 0$

proof –

from *assms* **obtain** δ **where** $\delta\text{-pos}: \delta > 0$

and $\forall y. |y - x| < \delta \longrightarrow |g\ y - g\ x| < -(g\ x)/2$

by (*metis continuous-at-eps-delta half-gt-zero neg-0-less-iff-less*)

then have $\forall y. |y - x| < \delta \longrightarrow -g\ y > 0$

by (*smt (verit, best) field-sum-of-halves*)

then show *?thesis*

using $\delta\text{-pos}$ *neg-0-less-iff-less* **by** *blast*

qed

lemma *cont-at-neg-imp-loc-neg'*:

fixes $g :: \text{real} \Rightarrow \text{real}$ **and** $x :: \text{real}$

assumes *continuous (at x) g* **and** $g\ x < 0$

shows $\exists \Delta > 0. \forall \delta. 0 < \delta \wedge \delta \leq \Delta \longrightarrow (\forall y. |y - x| < \delta \longrightarrow g\ y < 0)$

proof –

from *assms* **obtain** δ **where** $\delta\text{-pos}: \delta > 0$

and $H: \forall y. |y - x| < \delta \longrightarrow -(g\ y) > 0$

by (*smt (verit) cont-at-neg-imp-loc-neg*)

have $\forall \delta' \leq \delta. \forall y. |y - x| < \delta' \longrightarrow -(g\ y) > 0$

proof *clarify*

fix $\delta'\ y :: \text{real}$

assume $\delta' \leq \delta$ **and** $|y - x| < \delta'$

then show $-(g\ y) > 0$

using H **by** *auto*

qed

then show *?thesis*

using $\delta\text{-pos}$ *neg-0-less-iff-less* **by** *blast*

qed

end

2 Minimizers in Topological and Metric Spaces

theory *Minimizers-Definition*

imports *Auxiliary-Facts*

begin

2.1 Abstract Topological Definitions

definition *global-minimizer* :: $(\text{'a}::\text{topological-space} \Rightarrow \text{real}) \Rightarrow \text{'a} \Rightarrow \text{bool}$ **where**
global-minimizer $f\ x\text{-star} \longleftrightarrow (\forall x. f\ x\text{-star} \leq f\ x)$

definition *local-minimizer-on* :: $(\text{'a}::\text{topological-space} \Rightarrow \text{real}) \Rightarrow \text{'a} \Rightarrow \text{'a set} \Rightarrow$

bool **where**

local-minimizer-on f x -*star* $U \longleftrightarrow (\text{open } U \wedge x\text{-star} \in U \wedge (\forall x \in U. f\ x\text{-star} \leq f\ x))$

definition *local-minimizer* $:: ('a::\text{topological-space} \Rightarrow \text{real}) \Rightarrow 'a \Rightarrow \text{bool}$ **where**

local-minimizer f x -*star* $\longleftrightarrow (\exists U. \text{open } U \wedge x\text{-star} \in U \wedge (\forall x \in U. f\ x\text{-star} \leq f\ x))$

definition *isolated-local-minimizer-on* $:: ('a::\text{topological-space} \Rightarrow \text{real}) \Rightarrow 'a \Rightarrow 'a$
set $\Rightarrow \text{bool}$ **where**

isolated-local-minimizer-on f x -*star* $U \longleftrightarrow$
 $(\text{local-minimizer-on } f\ x\text{-star } U \wedge (\{x \in U. \text{local-minimizer } f\ x\} = \{x\text{-star}\}))$

definition *isolated-local-minimizer* $:: ('a::\text{topological-space} \Rightarrow \text{real}) \Rightarrow 'a \Rightarrow \text{bool}$
where

isolated-local-minimizer f x -*star* \longleftrightarrow
 $(\exists U. \text{local-minimizer-on } f\ x\text{-star } U \wedge (\{x \in U. \text{local-minimizer } f\ x\} = \{x\text{-star}\}))$

definition *strict-local-minimizer-on* $:: ('a::\text{topological-space} \Rightarrow \text{real}) \Rightarrow 'a \Rightarrow 'a$
set $\Rightarrow \text{bool}$ **where**

strict-local-minimizer-on f x -*star* $U \longleftrightarrow$
 $(\text{open } U \wedge x\text{-star} \in U \wedge (\forall x \in U - \{x\text{-star}\}. f\ x\text{-star} < f\ x))$

definition *strict-local-minimizer* $:: ('a::\text{topological-space} \Rightarrow \text{real}) \Rightarrow 'a \Rightarrow \text{bool}$
where

strict-local-minimizer f x -*star* $\longleftrightarrow (\exists U. \text{strict-local-minimizer-on } f\ x\text{-star } U)$

2.2 Metric Space Reformulations

lemma *local-minimizer-on-def2*:

fixes $f :: 'a::\text{metric-space} \Rightarrow \text{real}$

assumes *local-minimizer* f x -*star*

shows $\exists N > 0. \forall x \in \text{ball } x\text{-star } N. f\ x\text{-star} \leq f\ x$

proof –

from *assms* **obtain** U **where**

open U x -*star* $\in U$ **and** *local-min*: $\forall x \in U. f\ x\text{-star} \leq f\ x$

unfolding *local-minimizer-def* **by** *auto*

then obtain N **where** *N-pos*: $N > 0$ **and** *ball-in-U*: $\text{ball } x\text{-star } N \subseteq U$

using *open-contains-ball* **by** *blast*

hence $\forall x \in \text{ball } x\text{-star } N. f\ x\text{-star} \leq f\ x$

using *ball-in-U local-min* **by** *auto*

thus *?thesis*

using *N-pos* **by** *auto*

qed

lemma *local-minimizer-def2*:

fixes $f :: 'a::\text{metric-space} \Rightarrow \text{real}$

assumes *local-minimizer* f x -*star*

shows $\exists N > 0. \forall x. \text{dist } x\ x\text{-star} < N \longrightarrow f\ x\text{-star} \leq f\ x$

proof –
from *assms* **obtain** U **where**
open U $x\text{-star} \in U$ **and** *local-min*: $\forall x \in U. f\ x\text{-star} \leq f\ x$
unfolding *local-minimizer-def* **by** *auto*
then obtain N **where** *N-pos*: $N > 0$ **and** *ball-in-U*: $\text{ball}\ x\text{-star}\ N \subseteq U$
using *open-contains-ball* **by** *blast*
hence $\forall x. \text{dist}\ x\ x\text{-star} < N \longrightarrow x \in \text{ball}\ x\text{-star}\ N$
by (*subst mem-ball, simp add: dist-commute*)
hence $\forall x. \text{dist}\ x\ x\text{-star} < N \longrightarrow f\ x\text{-star} \leq f\ x$
using *ball-in-U local-min* **by** *blast*
thus *?thesis*
using *N-pos* **by** *auto*
qed

lemma *isolated-local-minimizer-on-def2*:
fixes $f :: 'a::\text{metric-space} \Rightarrow \text{real}$
assumes *isolated-local-minimizer-on* $f\ x\text{-star}\ U$
shows $\exists N > 0. \forall x \in \text{ball}\ x\text{-star}\ N. (\text{local-minimizer}\ f\ x \longrightarrow x = x\text{-star})$
proof –
from *assms* **have**
local-minimizer-on $f\ x\text{-star}\ U$
and *unique-min*: $\{x \in U. \text{local-minimizer}\ f\ x\} = \{x\text{-star}\}$
unfolding *isolated-local-minimizer-on-def* **by** *auto*
then obtain N **where** *N-pos*: $N > 0$ **and** *ball-in-U*: $\text{ball}\ x\text{-star}\ N \subseteq U$
using *open-contains-ball* **by** (*metis local-minimizer-on-def*)
have $\forall x \in \text{ball}\ x\text{-star}\ N. \text{local-minimizer}\ f\ x \longrightarrow x = x\text{-star}$
proof (*clarify*)
fix x
assume $x \in \text{ball}\ x\text{-star}\ N$
then have $x \in U$ **using** *ball-in-U* **by** *auto*
moreover assume *local-minimizer* $f\ x$
hence $x \in \{x \in U. \text{local-minimizer}\ f\ x\}$ **using** $\langle x \in U \rangle$ **by** *auto*
hence $x \in \{x\text{-star}\}$ **using** *unique-min* **by** *auto*
ultimately show $x = x\text{-star}$
by *simp*
qed
thus *?thesis* **using** *N-pos* **by** *auto*
qed

lemma *isolated-local-minimizer-def2*:
fixes $f :: 'a::\text{metric-space} \Rightarrow \text{real}$
assumes *isolated-local-minimizer* $f\ x\text{-star}$
shows $\exists N > 0. \forall x \in \text{ball}\ x\text{-star}\ N. (\text{local-minimizer}\ f\ x \longrightarrow x = x\text{-star})$
proof –
from *assms* **obtain** U **where**
local-minimizer-on $f\ x\text{-star}\ U$
and *unique-min*: $\{x \in U. \text{local-minimizer}\ f\ x\} = \{x\text{-star}\}$
unfolding *isolated-local-minimizer-def* **by** *auto*
then obtain N **where** *N-pos*: $N > 0$ **and** *ball-in-U*: $\text{ball}\ x\text{-star}\ N \subseteq U$

```

    using open-contains-ball by (metis local-minimizer-on-def)
  have  $\forall x \in \text{ball } x\text{-star } N. \text{local-minimizer } f x \longrightarrow x = x\text{-star}$ 
  proof (clarify)
    fix x
    assume  $x \in \text{ball } x\text{-star } N$ 
    then have  $x \in U$  using ball-in-U by auto
    moreover assume local-minimizer f x
    hence  $x \in \{x \in U. \text{local-minimizer } f x\}$  using  $\langle x \in U \rangle$  by auto
    hence  $x \in \{x\text{-star}\}$  using unique-min by auto
    ultimately show  $x = x\text{-star}$  by simp
  qed
  thus ?thesis using N-pos by auto
qed

```

```

lemma strict-local-minimizer-on-def2:
  fixes f :: 'a::metric-space  $\Rightarrow$  real
  assumes strict-local-minimizer-on f x-star U
  shows  $\exists N > 0. \forall x \in \text{ball } x\text{-star } N - \{x\text{-star}\}. f x\text{-star} < f x$ 
  proof -
    from assms have
      open U x-star  $\in U$  and strict-min:  $\forall x \in U - \{x\text{-star}\}. f x\text{-star} < f x$ 
    unfolding strict-local-minimizer-on-def by auto
    then obtain N where N-pos:  $N > 0$  and ball-in-U:  $\text{ball } x\text{-star } N \subseteq U$ 
      using open-contains-ball by metis
    have  $\forall x \in \text{ball } x\text{-star } N - \{x\text{-star}\}. f x\text{-star} < f x$ 
    proof
      fix x
      assume  $x \in \text{ball } x\text{-star } N - \{x\text{-star}\}$ 
      hence  $x \in U - \{x\text{-star}\}$  using ball-in-U by auto
      thus  $f x\text{-star} < f x$ 
        using strict-min by auto
    qed
    thus ?thesis using N-pos by auto
  qed

```

```

lemma strict-local-minimizer-def2:
  fixes f :: 'a::metric-space  $\Rightarrow$  real
  assumes strict-local-minimizer f x-star
  shows  $\exists N > 0. \forall x \in \text{ball } x\text{-star } N - \{x\text{-star}\}. f x\text{-star} < f x$ 
  proof -
    from assms obtain U where
      strict-local-minimizer-on f x-star U
    unfolding strict-local-minimizer-def by auto
    then have
      open U x-star  $\in U$  and strict-min:  $\forall x \in U - \{x\text{-star}\}. f x\text{-star} < f x$ 
    unfolding strict-local-minimizer-on-def by auto
    then obtain N where N-pos:  $N > 0$  and ball-in-U:  $\text{ball } x\text{-star } N \subseteq U$ 
      using open-contains-ball by metis
    have  $\forall x \in \text{ball } x\text{-star } N - \{x\text{-star}\}. f x\text{-star} < f x$ 

```

```

proof
  fix  $x$ 
  assume  $x \in \text{ball } x\text{-star } N - \{x\text{-star}\}$ 
  hence  $x \in U - \{x\text{-star}\}$  using ball-in-U by auto
  thus  $f\ x\text{-star} < f\ x$ 
    using strict-min by auto
  qed
  thus ?thesis using N-pos by auto
qed

lemma local-minimizer-neighborhood:
  fixes  $f :: \text{real} \Rightarrow \text{real}$ 
  assumes loc-min: local-minimizer  $f\ x\text{-min}$ 
  shows  $\exists \delta > 0. \forall h. |h| < \delta \longrightarrow f\ (x\text{-min} + h) \geq f\ x\text{-min}$ 
proof -
  obtain  $N$  where N-pos:  $N > 0$  and N-prop:  $\forall x. \text{dist } x\ x\text{-min} < N \longrightarrow f\ x\text{-min} \leq f\ x$ 
    using local-minimizer-def2[OF loc-min] by auto
  then have  $\forall h. \text{abs } h < N \longrightarrow f\ (x\text{-min} + h) \geq f\ x\text{-min}$ 
    by (simp add: dist-real-def)
  then show ?thesis
    using N-pos by blast
qed

lemma local-minimizer-from-neighborhood:
  fixes  $f :: \text{real} \Rightarrow \text{real}$  and  $x\text{-min} :: \text{real}$ 
  assumes  $\exists \delta > 0. \forall x. |x - x\text{-min}| < \delta \longrightarrow f\ x\text{-min} \leq f\ x$ 
  shows local-minimizer  $f\ x\text{-min}$ 
proof -
  from assms obtain  $\delta$  where delta-pos:  $\delta > 0$  and H:  $\forall x. |x - x\text{-min}| < \delta \longrightarrow f\ x\text{-min} \leq f\ x$ 
    by auto
  obtain  $U$  where U-def:  $U = \{x. |x - x\text{-min}| < \delta\}$ 
    by simp
  then have open  $U$ 
    by (smt (verit) dist-commute dist-real-def mem-Collect-eq metric-space-class.open-ball subsetI topological-space-class.openI)
  moreover have  $x\text{-min} \in U$ 
    using U-def delta-pos by force
  moreover have  $\forall x \in U. f\ x\text{-min} \leq f\ x$ 
    using H U-def by blast
  ultimately show ?thesis
    unfolding local-minimizer-def by auto
qed

end

```

3 Minimizer Implications

```
theory First-Order-Conditions
  imports Minimizers-Definition
begin
```

```
notation norm (||-||)
```

3.1 Implications for a Given Minimizer Type

lemma *strict-local-minimizer-imp-local-minimizer*:

assumes *strict-local-minimizer f x-star*

shows *local-minimizer f x-star*

by (*smt (verit) Diff-iff assms local-minimizer-def singletonD strict-local-minimizer-def strict-local-minimizer-on-def*)

lemma *isolated-local-minimizer-imp-strict*:

assumes *isolated-local-minimizer f x-star*

shows *strict-local-minimizer f x-star*

proof –

– From *isolated_local_minimizer* we obtain an open set U such that x^* is the *only* local minimizer.

from *assms* **obtain** U **where** *iso-props*:

isolated-local-minimizer-on f x-star U

unfolding *isolated-local-minimizer-def*

using *isolated-local-minimizer-on-def* **by** *blast*

– Unpack *isolated_local_minimizer_on*: x^* is a *local_minimizer_on* U , and x^* is unique.

from *iso-props* **have** *lm-on: local-minimizer-on f x-star U*

unfolding *isolated-local-minimizer-on-def* **using** *local-minimizer-on-def* **by** *presburger*

moreover from *iso-props* **have** *unique-min: $\{x \in U. \text{local-minimizer } f \ x\} = \{x\text{-star}\}$*

unfolding *isolated-local-minimizer-on-def* **by** *auto*

– From *local_minimizer_on*, we have: U open, $x^* \in U$, and $\forall x \in U. f(x^*) \leq f(x)$.

from *lm-on* **have** *open-U: open U* **and** *x-in-U: x-star \in U* **and** *le-prop: $\forall x \in U. f \ x\text{-star} \leq f \ x$*

unfolding *local-minimizer-on-def* **by** *auto*

– Assume, for contradiction, that x^* is not a strict local minimizer. Then there exists $y \in U \setminus \{x^*\}$ with $f(y) \leq f(x^*)$.

show *strict-local-minimizer f x-star*

proof (*rule ccontr*)

assume \neg *strict-local-minimizer f x-star*

then obtain y where y -props:

$y \in U - \{x\text{-star}\}$ **and** $f y \leq f x\text{-star}$

unfolding *strict-local-minimizer-def strict-local-minimizer-on-def*

by (*smt (verit, ccfv-SIG) open-U x-in-U*)

from y -props have $y \in U$ and $y \neq x\text{-star}$

by *auto*

— We already have $f(x^*) \leq f(y)$ from $\forall x \in U. f x\text{-star} \leq f x$ and $y \in U$. Together with $f(y) \leq f(x^*)$, this yields $f(x^*) = f(y)$.

from *le-prop* $\langle y \in U \rangle$ **have** $f x\text{-star} \leq f y$

by *auto*

with $\langle f y \leq f x\text{-star} \rangle$ **have** $f x\text{-star} = f y$

by *auto*

— Now we show that y is also a local minimizer, contradicting the uniqueness of x^* . To prove this, we must exhibit an open set V around y such that $f(y) \leq f(x)$ for all $x \in V$.

have *local-minimizer* $f y$

proof —

— Since U is open and $y \in U$, there exists an open set $V \subseteq U$ containing y .

obtain V **where** *open* V **and** $y \in V$ **and** $V \subseteq U$

using $\langle \text{open } U \rangle \langle y \in U \rangle$ *open-subset* **by** *auto*

— On this subset, $f(y) = f(x^*) \leq f(x)$ for all $x \in V$ (since $V \subseteq U$).

moreover from *le-prop* **and** $\langle f x\text{-star} = f y \rangle$ **have** $\forall x \in V. f y \leq f x$

using *calculation(3)* **by** *auto*

ultimately show *local-minimizer* $f y$

unfolding *local-minimizer-def local-minimizer-on-def* **by** *auto*

qed

— Since y is a local minimizer and $y \in U$, we have $y \in \{x \in U. \text{local_minimizer } f x\}$. By uniqueness, $\{x \in U. \text{local_minimizer } f x\} = \{x^*\}$, hence $y = x^*$, contradicting $y \neq x^*$.

hence $y \in \{x \in U. \text{local_minimizer } f x\}$

by (*simp add: $\langle y \in U \rangle$*)

with *unique-min* **have** $y = x\text{-star}$ **by** *auto*

thus *False* **using** $\langle y \neq x\text{-star} \rangle$ **by** *contradiction*

qed

— Having reached a contradiction under the assumption that x^* is not a strict local minimizer, it follows that x^* must indeed be a strict local minimizer.

qed

3.2 Characterization of Non-Isolated Minimizers

lemma *not-isolated-minimizer-def*:

assumes *local-minimizer f x-star*

shows $(\exists x\text{-seq}::\text{nat} \Rightarrow \text{real}. (\forall n. \text{local-minimizer } f (x\text{-seq } n) \wedge x\text{-seq } n \neq x\text{-star})) \wedge ((x\text{-seq} \longrightarrow x\text{-star}) \text{ at-top}) = (\neg \text{isolated-local-minimizer } f x\text{-star})$

proof(*safe*)

show $\bigwedge x\text{-seq}. \text{isolated-local-minimizer } f x\text{-star} \Longrightarrow \forall n. \text{local-minimizer } f (x\text{-seq } n) \wedge x\text{-seq } n \neq x\text{-star} \Longrightarrow x\text{-seq} \longrightarrow x\text{-star} \Longrightarrow \text{False}$

proof –

fix $x\text{-seq} :: \text{nat} \Rightarrow \text{real}$

assume *x-star-isolated-minimizer: isolated-local-minimizer f x-star*

assume *with-sequence-of-local-minimiziers: $\forall n. \text{local-minimizer } f (x\text{-seq } n) \wedge x\text{-seq } n \neq x\text{-star}$*

assume *converging-to-x-star: $x\text{-seq} \longrightarrow x\text{-star}$*

have *open-ball-with-unique-min: $\exists N > 0. \forall x \in \text{ball } x\text{-star } N. (\text{local-minimizer } f x \longrightarrow x = x\text{-star})$*

by (*simp add: isolated-local-minimizer-def2 x-star-isolated-minimizer*)

then obtain N **where** *N-pos: $N > 0$ and N-prop: $\forall x \in \text{ball } x\text{-star } N. (\text{local-minimizer } f x \longrightarrow x = x\text{-star})$*

by *blast*

– Use convergence to show x_{seq} eventually lies in $\text{ball}(x^*, N)$.

from *converging-to-x-star* **have** $\exists M. \forall n \geq M. x\text{-seq } n \in \text{ball } x\text{-star } N$

by (*metis LIMSEQ-iff-nz N-pos dist-commute mem-ball*)

then obtain M **where** *M-def: $\forall n \geq M. x\text{-seq } n \in \text{ball } x\text{-star } N$*

by *auto*

then show *False*

by (*meson N-prop linorder-not-le order-less-irrefl with-sequence-of-local-minimiziers*)

qed

next

show $\neg \text{isolated-local-minimizer } f x\text{-star} \Longrightarrow \exists x\text{-seq}. (\forall n. \text{local-minimizer } f (x\text{-seq } n) \wedge x\text{-seq } n \neq x\text{-star}) \wedge x\text{-seq} \longrightarrow x\text{-star}$

proof(*rule ccontr*)

assume *not-isolated-minimizer: $\neg \text{isolated-local-minimizer } f x\text{-star}$*

assume *BWOC: $\nexists x\text{-seq}. (\forall n. \text{local-minimizer } f (x\text{-seq } n) \wedge x\text{-seq } n \neq x\text{-star}) \wedge x\text{-seq} \longrightarrow x\text{-star}$*

have $\exists N > 0. \forall x. \text{dist } x x\text{-star} < N \longrightarrow f x\text{-star} \leq f x$

by (*simp add: assms local-minimizer-def2*)

then obtain N **where** *N-pos: $(N::\text{nat}) > 0$ and x-star-min-on-N-ball: $\forall x. \text{dist } x x\text{-star} < 1/\text{real } N \longrightarrow f x\text{-star} \leq f x$*

by (*metis dual-order.strict-trans ex-inverse-of-nat-less inverse-eq-divide*)

obtain $S\text{-n} :: \text{nat} \Rightarrow \text{real}$ **set where** *S-n-def: $S\text{-n} = (\lambda n. \{x. \text{dist } x x\text{-star} < 1/\text{real } n + N\})$*

by *blast*

from *not-isolated-minimizer*

have *non-isolated: $\forall U. \text{local-minimizer-on } f x\text{-star } U \longrightarrow (\exists y \in U. y \neq x\text{-star} \wedge \text{local-minimizer } f y)$*

by (smt (verit, best) Collect-cong assms isolated-local-minimizer-def local-minimizer-on-def singleton-conv2)

have $\forall n::nat. \exists x. x \in S-n$

proof (intro allI)

fix $n::nat$

have pos-radius: $1 / (\text{real } n + N) > 0$

using N-pos by simp

obtain U where U-def: $U = \text{ball } x\text{-star } (1 / (\text{real } n + N))$ and open-U: open U and U-contains-x-star: $x\text{-star} \in U$

using pos-radius by auto

have U-contained-in-Inverse-N-Ball: $\forall x \in U. \text{dist } x \text{ } x\text{-star} < 1 / N$

proof (safe)

fix $x::\text{real}$

assume x-in-U: $x \in U$

then have $\text{dist } x \text{ } x\text{-star} < (1 / (\text{real } n + N))$

by (simp add: U-def dist-commute)

also have $\dots \leq 1 / \text{real } N$

by (simp add: N-pos frac-le)

finally show $\text{dist } x \text{ } x\text{-star} < 1 / \text{real } N$.

qed

have ball-non-empty: $\exists y \in U. y \neq x\text{-star} \wedge \text{local-minimizer } f y$

proof -

have local-minimizer-on f x-star U

by (simp add: U-contains-x-star U-contained-in-Inverse-N-Ball local-minimizer-on-def open-U x-star-min-on-N-ball)

then show $\exists y \in U. y \neq x\text{-star} \wedge \text{local-minimizer } f y$

by (simp add: non-isolated)

qed

then obtain y where y-in-ball: $y \in U$ and $y \neq x\text{-star}$ and local-minimizer f y

by blast

then show $\exists x. x \in S-n$

by (smt (verit, best) S-n-def U-def dist-commute mem-Collect-eq mem-ball)

qed

then obtain x-seq where x-seq-def: $\forall n. x\text{-seq } n \in S-n$

by metis

have x-seq-converges-to-x-star: $x\text{-seq} \longrightarrow x\text{-star}$

proof (rule LIMSEQ-I)

fix $r::\text{real}$

assume r-pos: $0 < r$

obtain n-min where n-min-def: $1 / (\text{real } n\text{-min} + N) < r$

using real-arch-inverse N-pos r-pos

by (smt (verit, ccfv-SIG) frac-le inverse-eq-divide inverse-positive-iff-positive)

show $\exists no. \forall n \geq no. \text{norm } (x\text{-seq } n - x\text{-star}) < r$

proof (intro exI allI impI)

```

fix n
assume n ≥ n-min
then have n-large-enough: 1 / (real n + N) ≤ 1 / (real n-min + N)
  using N-pos by (subst frac-le, simp-all)
have dist (x-seq n) x-star < 1 / (real n + N)
  using x-seq-def S-n-def by auto
also have ... ≤ 1 / (real n-min + N)
  using n-large-enough by auto
also have ... < r
  using n-min-def by auto
finally show norm (x-seq n - x-star) < r
  by (simp add: dist-real-def)
qed
qed
have ∃ x-seq. (∀ n. local-minimizer f (x-seq n) ∧ x-seq n ≠ x-star) ∧ x-seq
  → x-star
  using S-n-def x-seq-converges-to-x-star x-seq-def by blast
then show False
  using BWOC by auto
qed
qed

```

3.3 First-Order Condition

theorem *Fermat's-theorem-on-stationary-points:*

```

fixes f :: real ⇒ real
assumes (f has-derivative f') (at x-min)
assumes local-minimizer f x-min
shows (deriv f) x-min = 0
by (metis assms has-derivative-imp differential-zero-maxmin local-minimizer-def)

```

definition *stand-basis-vector* :: 'n::finite ⇒ realⁿ — the i-th standard basis vector

```

where stand-basis-vector i = (χ j. if j = i then 1 else 0)

```

lemma *stand-basis-vector-index[simp]:* (stand-basis-vector i) \$ j = (if j = i then (1::real) else 0)

```

by (simp add: stand-basis-vector-def)

```

lemma *stand-basis-vector-nonzero[simp]:* stand-basis-vector i ≠ 0

```

by (smt (verit, del-insts) stand-basis-vector-index zero-index)

```

lemma *norm-stand-basis-vector[simp]:* norm (stand-basis-vector i) = 1

```

by (smt (verit, best) axis-nth component-le-norm-cart norm-axis-1 norm-le-componentwise-cart
  real-norm-def stand-basis-vector-index)

```

lemma *inner-stand-basis-vector[simp]:* inner (stand-basis-vector i) (stand-basis-vector j) = (if i = j then 1 else 0)

```

by (metis axis-nth cart-eq-inner-axis norm-eq-1 norm-stand-basis-vector stand-basis-vector-index)

```

vector-eq)

lemma *Basis-characterisation:*

stand-basis-vector $i \in (\text{Basis} :: (\text{real}^n \text{ set}) \text{ and}$
 $\forall b \in (\text{Basis} :: (\text{real}^n \text{ set}). \exists i. b = \text{stand-basis-vector } i$
by (*metis* (*no-types*, *lifting*) *Basis-real-def axis-in-Basis-iff cart-eq-inner-axis*
inner-stand-basis-vector insert-iff norm-axis-1 norm-eq-1 stand-basis-vector-index
vector-eq,
metis axis-index axis-nth cart-eq-inner-axis inner-stand-basis-vector stand-basis-vector-index
vector-eq)

lemma *stand-basis-expansion:*

fixes $x :: \text{real}^n$
shows $x = (\sum_{j \in \text{UNIV}} (x \ \$ \ j) *_{\mathbb{R}} \text{stand-basis-vector } j)$
proof –
have $(\sum_{j \in \text{UNIV}} (x \ \$ \ j) *_{\mathbb{R}} \text{stand-basis-vector } j) \ \$ \ k = x \ \$ \ k$ **for** k
proof –
have $(\sum_{j \in \text{UNIV}} (x \ \$ \ j) *_{\mathbb{R}} \text{stand-basis-vector } j) \ \$ \ k$
 $= (\sum_{j \in \text{UNIV}} (x \ \$ \ j) * (\text{stand-basis-vector } j \ \$ \ k))$
by *simp*
also have $\dots = (\sum_{j \in \text{UNIV}} (x \ \$ \ j) * (\text{if } j = k \text{ then } 1 \text{ else } 0))$
by (*smt* (*verit*, *best*) *stand-basis-vector-index sum.cong*)
also have $\dots = (\sum_{j \in \text{UNIV}} (\text{if } j = k \text{ then } x \ \$ \ j \text{ else } 0))$
by (*smt* (*verit*, *best*) *mult-cancel-left1 mult-cancel-right1 sum.cong*)
also have $\dots = x \ \$ \ k$
by (*subst sum.delta*, *simp-all*)
finally show *?thesis*.
qed
thus *?thesis*
by (*simp add: vec-eq-iff*)
qed

lemma *has-derivative-affine:*

fixes $a \ v :: 'a :: \text{real-normed-vector}$
shows $((\lambda t. a + t *_{\mathbb{R}} v) \text{ has-derivative } (\lambda h. h *_{\mathbb{R}} v)) \text{ (at } x)$
unfolding *has-derivative-def*
proof *safe*
have $a + y *_{\mathbb{R}} v - (a + \text{netlimit } (\text{at } x) *_{\mathbb{R}} v) - (y - \text{netlimit } (\text{at } x)) *_{\mathbb{R}} v = 0$
if $y \neq \text{netlimit } (\text{at } x)$ **for** y
by (*simp add: cross3-simps(32)*)
then show $(\lambda y. (a + y *_{\mathbb{R}} v - (a + \text{netlimit } (\text{at } x) *_{\mathbb{R}} v) - (y - \text{netlimit } (\text{at } x)) *_{\mathbb{R}} v) /_{\mathbb{R}} \|y - \text{netlimit } (\text{at } x)\|) -x \rightarrow 0$
by (*simp add: scaleR-left-diff-distrib*)
show *bounded-linear* $(\lambda h. h *_{\mathbb{R}} v)$
by (*simp add: bounded-linearI' vector-space-assms(2)*)
qed

theorem *Fermat's-theorem-on-stationary-points-mult:*

fixes $f :: \text{real}^n \Rightarrow \text{real}$

assumes *der-f*: (*f* has-derivative *f'*) (at *x-min*)
assumes *min-f*: local-minimizer *f* *x-min*
shows *GDERIV* *f* *x-min* :> 0
proof –
– Show that *f'* kills every standard-basis vector.

{
 fix *i* :: 'n
 – Define the 1D slice $g_i(t) = f(x_{\min} + t \cdot e_i)$.
 let *?g* = $\lambda t :: \text{real. } f(x_{\min} + t *_{\mathbb{R}} \text{stand-basis-vector } i)$

 – Chain rule gives $g'_i(0) = f'(e_i)$.
 from *has-derivative-affine* **have** *g-der*:
 ($\lambda t. f(x_{\min} + t *_{\mathbb{R}} \text{stand-basis-vector } i)$
 has-derivative ($\lambda h. f'(h *_{\mathbb{R}} \text{stand-basis-vector } i)$) (at 0))
 by (*metis* (*no-types*) *arithmetic-simps*(50) *der-f* *has-derivative-compose* *scaleR-simps*(1))

 – 0 is a local minimizer of g_i because x_{\min} is one for *f*.
 have *g-min*: local-minimizer *?g* 0
 proof(*rule* *local-minimizer-from-neighborhood*)
 obtain δ **where** *delta-pos*: $\delta > 0$
 and *mono*: $\bigwedge x. \text{dist } x_{\min} x < \delta \implies f x \geq f x_{\min}$
 by (*metis* *assms*(2) *dist-commute* *local-minimizer-def2*)

 have $\forall x. |x - 0| < \delta \implies f(x_{\min} + 0 *_{\mathbb{R}} \text{stand-basis-vector } i) \leq f(x_{\min} + x *_{\mathbb{R}} \text{stand-basis-vector } i)$
 using *mono* **by** (*simp* *add*: *dist-norm*)
 then show $\exists \delta > 0. \forall x. |x - 0| < \delta \implies f(x_{\min} + 0 *_{\mathbb{R}} \text{stand-basis-vector } i) \leq f(x_{\min} + x *_{\mathbb{R}} \text{stand-basis-vector } i)$
 using *delta-pos* **by** *blast*
 qed

 – Apply the 1-D Fermat lemma to g_i .
 from *Fermat's-theorem-on-stationary-points*
 have $f'(\text{stand-basis-vector } i) = 0$
 using *g-der* *g-min* **by** (*metis* *has-derivative-imp* *scale-one*)
}

– Collecting the result for every *i*:
hence *zero-on-basis*: $\bigwedge i. f'(\text{stand-basis-vector } i) = 0$.

– Use linearity and the coordinate expansion to show $f' = 0$ everywhere.

{
 fix *v* :: real^n
 – Expand $v = \sum_j v_j \cdot e_j$ and push *f'* through the finite sum.
 have $f' v = 0$
 proof –
 have $f' v = f'(\sum_{j \in \text{UNIV.}} (v \$ j) *_{\mathbb{R}} \text{stand-basis-vector } j)$
 by (*metis* *stand-basis-expansion*)
}

```

    also have ... = ( $\sum_{j \in UNIV}. (v \ \$ \ j) *_{\mathbb{R}} f' (stand-basis-vector \ j)$ )
    by (smt (verit) assms differential-zero-maxmin local-minimizer-def scale-eq-0-iff
sum.neutral)
    also have ... = 0
    using zero-on-basis by simp
    finally show ?thesis.
  qed
}
hence f'-zero: f' = ( $\lambda \cdot 0$ )
  by (simp add: fun-eq-iff)

— Translate  $f' = 0$  into the gradient statement.
have (f has-derivative ( $\lambda h. 0$ )) (at x-min)
  using der-f f'-zero by simp
hence GDERIV f x-min :> ( $0 :: real^n$ )
  by (simp add: gderiv-def)
thus ?thesis.
qed

end

```

4 Second-Order Conditions

```

theory Second-Derivative-Test
  imports First-Order-Conditions
begin

```

4.1 Necessary Condition

```

lemma snd-derivative-nonneg-at-local-min-necessary:
  fixes f :: real  $\Rightarrow$  real
  assumes C2-cont-diff-at-xmin: C-k-on 2 f (U :: real set)
  assumes min-in-U: (x-min :: real)  $\in$  U
  assumes loc-min: local-minimizer f x-min
  shows deriv (deriv f) x-min  $\geq 0$ 
proof —
  have ( $\exists \ \varepsilon. 0 < \varepsilon \wedge \{x-min - \varepsilon .. x-min + \varepsilon\} \subset U$ )
  proof —
    have ( $\exists \ \varepsilon. 0 < \varepsilon \wedge ball \ x-min \ \varepsilon \subset U$ )
    by (smt C2-cont-diff-at-xmin C-k-on-def assms(2) ball-subset-cball cball-eq-ball-iff

        open-contains-cball-eq order-le-less-trans psubsetI)
    then show ?thesis
    by (metis Elementary-Metric-Spaces.open-ball cball-eq-atLeastAtMost cen-
tre-in-ball
        open-contains-cball order-trans-rules(21))
  qed
  then obtain  $\varepsilon$  where  $\varepsilon$ -pos:  $0 < \varepsilon$  and  $\varepsilon$ -def:  $\{x-min - \varepsilon .. x-min + \varepsilon\} \subset U$ 
  by blast

```

have $f\text{-diff}$: $(\forall y \in U. (f \text{ has-real-derivative } (\text{deriv } f) y) (at y))$
using $C2\text{-cont-diff } C2\text{-cont-diff-at-xmin}$ **by** $blast$
have $f'\text{-diff}$: $(\forall y \in U. (\text{deriv } f \text{ has-real-derivative } (\text{deriv } (\text{deriv } f)) y) (at y))$
using $C2\text{-cont-diff } C2\text{-cont-diff-at-xmin}$ **by** $blast$
have $f''\text{-contin}$: $continuous\text{-on } U (\text{deriv } (\text{deriv } f))$
using $C2\text{-cont-diff } assms(1)$ **by** $blast$

have $f'\text{-0}$: $(\text{deriv } f) x\text{-min} = 0$
using $Fermat's\text{-theorem-on-stationary-points}$
by $(meson\ assms(2,3) f\text{-diff has-field-derivative-imp-has-derivative})$

— By local minimality at x_{\min} , there is a $\delta > 0$ such that for all h with $|h| < \delta$, we have $f(x_{\min} + h) \geq f(x_{\min})$.
obtain δ **where** $\delta\text{-pos}$: $\delta > 0$ **and** $\delta\text{-prop}$: $\forall h. |h| < \delta \longrightarrow f(x\text{-min} + h) \geq f\ x\text{-min}$
by $(meson\ assms(3) local\ minimizer\ neighborhood)$

from $f'\text{-0}$ **have** $second\ deriv\ limit\ at\ x\text{-min}$:
 $((\lambda h. (\text{deriv } f (x\text{-min} + h)) / h) \longrightarrow \text{deriv } (\text{deriv } f) x\text{-min}) (at 0)$
by $(smt (verit, best) DERIV\text{-def } Lim\text{-cong-within } assms(2) f'\text{-diff})$

show $?thesis$
proof $(rule\ ccontr)$
assume $\neg 0 \leq \text{deriv } (\text{deriv } f) x\text{-min}$
then have $BWOC$: $0 > \text{deriv } (\text{deriv } f) x\text{-min}$
by $auto$
then obtain Δ **where** $\Delta\text{-pos}$: $\Delta > 0$ **and**
 $\Delta\text{-def}$: $\forall \delta. 0 < \delta \wedge \delta \leq \Delta \longrightarrow (\forall y. |y - x\text{-min}| < \delta \longrightarrow \text{deriv } (\text{deriv } f) y < 0)$
by $(metis\ C2\text{-cont-diff-at-xmin } C\text{-k-on-def } min\text{-in-}U\ at\text{-within-open } cont\text{-at-neg-imp-loc-neg}'\ continuous\text{-on-eq-continuous-within } f''\text{-contin})$

— Choose h with $0 < h < \min\{\delta, \Delta\}$ so that $x_{\min} + h \in U$.
obtain h **where** $h\text{-def}$: $h = \min \varepsilon (min (\delta/2) \Delta)$ **and** $h\text{-pos}$: $0 < h$
using $\varepsilon\text{-pos } \delta\text{-pos } \Delta\text{-pos}$ **by** $fastforce$
have $h\text{-lt}$: $h \leq \varepsilon \wedge h < \delta \wedge h \leq \Delta$
using $\delta\text{-pos } h\text{-def}$ **by** $linarith$
have $neigh\text{-in-}U$: $x\text{-min} + h \in \{x\text{-min} - \varepsilon .. x\text{-min} + \varepsilon\}$
using $h\text{-def } h\text{-pos}$ **by** $fastforce$

have $f(x\text{-min} + h) < f\ x\text{-min}$
proof $(rule\ DERIV\text{-neg-imp-decreasing-open}$ **[where** $a = x\text{-min}$ **and** $f = f$ **and** $b = x\text{-min} + h$ $)$
show $x\text{-min} < x\text{-min} + h$
using $h\text{-pos}$ **by** $simp$
next
have $\{x\text{-min}..x\text{-min} + h\} \subset U$
using $\varepsilon\text{-def } dual\text{-order.strict-trans2 } neigh\text{-in-}U$ **by** $auto$
then show $continuous\text{-on } \{x\text{-min}..x\text{-min} + h\} f$

by (*meson C2-cont-diff C2-cont-diff-at-xmin continuous-on-subset differentiable-imp-continuous-on le-less*)

next

show $\bigwedge x. [x\text{-min} < x; x < x\text{-min} + h] \implies \exists y. (f \text{ has-real-derivative } y) (at\ x) \wedge y < 0$

proof –

fix $x :: real$

assume $x\text{-min-}lt\text{-}x: x\text{-min} < x$

assume $x\text{-}lt\text{-}x\text{-min-}pls\text{-}h: x < x\text{-min} + h$

have $x\text{-min-}x\text{-subset}: \{x\text{-min} .. x\} \subseteq \{x\text{-min} - \varepsilon .. x\text{-min} + \varepsilon\}$

using *neigh-in-U x-lt-xmin-pls-h* **by** *auto*

— By the Mean Value Theorem applied to f' on $[x_{\min}, x]$, there exists some c with $x_{\min} < c < x$ such that:

have $\exists z > x\text{-min}. z < x \wedge deriv\ f\ (x) - deriv\ f\ x\text{-min} = (x - x\text{-min}) * deriv\ (deriv\ f)\ z$

proof(*rule MVT2*)

show $x\text{-min} < x$

using $x\text{-min-}lt\text{-}x$ **by** *auto*

next

fix $y :: real$

assume $x\text{-min-}leq\text{-}y: x\text{-min} \leq y$

assume $y\text{-}leq\text{-}x: y \leq x$

from $x\text{-min-}x\text{-subset}$ **have** $y \in U$

using $\varepsilon\text{-def atLeastAtMost-iff } x\text{-min-}leq\text{-}y\ y\text{-}leq\text{-}x$ **by** *blast*

then show (*deriv f has-real-derivative deriv (deriv f) y*) (*at y*)

using $f'\text{-diff}$ **by** *blast*

qed

then obtain z **where** $z\text{-gt-}x\text{-min}: z > x\text{-min}$ **and**

$z\text{-lt-}x: z < x$ **and**

$z\text{-def}: deriv\ f\ (x) - deriv\ f\ x\text{-min} = (x - x\text{-min}) * deriv\ (deriv\ f)\ z$

(deriv f) z

by *blast*

then have $mvt\text{-}f': deriv\ f\ (x) = (x - x\text{-min}) * deriv\ (deriv\ f)\ z$

by (*simp add: f'-0*)

then have $x\text{-diff-}x\text{-min-}pos: x - x\text{-min} > 0$

using $\langle x\text{-min} < x \rangle$ **by** *simp*

then have $left\text{-}bound\text{-}satisfied: |z - x\text{-min}| < x - x\text{-min}$

using $\langle x\text{-min} < z \rangle \langle z < x \rangle$ **by** *auto*

then have $x - x\text{-min} < h$

using $\langle x < x\text{-min} + h \rangle$ **by** *simp*

then have $|z - x\text{-min}| < h$

using $left\text{-}bound\text{-}satisfied$ **by** *fastforce*

then have $deriv\ (deriv\ f)\ z < 0$

using $\Delta\text{-def } h\text{-lt } h\text{-pos}$ **by** *blast*

then have $deriv\ f\ x < 0$

```

    by (metis x-diff-xmin-pos mult-f' mult-pos-neg)
  moreover have  $x \in U$ 
    using xmin-x-subset
  by (meson  $\varepsilon$ -def atLeastAtMost-iff dual-order.strict-iff-not
    subset-eq verit-comp-simplify(2) x-min-lt-x)
  ultimately show  $\exists y. (f \text{ has-real-derivative } y) (at\ x) \wedge y < 0$ 
    using f-diff by blast
  qed
  qed
  then show False
    by (smt (verit, best)  $\delta$ -prop h-lt h-pos)
  qed
  qed

```

4.2 Sufficient Condition

lemma *second-derivative-test*:

```

  fixes  $f :: real \Rightarrow real$  and  $a :: real$  and  $b :: real$  and  $x\text{-min} :: real$ 
  assumes valid-interval:  $a < b$ 
  assumes twice-continuously-differentiable:  $C\text{-k-on } 2\ f\ \{a <..< < b\}$ 
  assumes min-exists:  $x\text{-min} \in \{a <..< < b\}$ 
  assumes fst-deriv-req:  $(deriv\ f)\ x\text{-min} = 0$ 
  assumes snd-deriv-req:  $deriv\ (deriv\ f)\ x\text{-min} > 0$ 
  shows loc-min: local-minimizer  $f\ x\text{-min}$ 
  proof -
    from twice-continuously-differentiable
  have  $f''\text{-cont}$ : continuous-on  $\{a <..< < b\}\ (deriv\ (deriv\ f))$ 
    by (metis  $C\text{-k-on-def}\ Suc\text{-1}\ lessI\ nat.\text{simps}(2)\ \text{second-derivative-alt-def}$ )
  then obtain  $\Delta$  where  $\Delta\text{-pos}$ :  $\Delta > 0$ 
    and  $\Delta\text{-prop}$ :  $\forall \delta. 0 < \delta \wedge \delta \leq \Delta \longrightarrow (\forall y. |y - x\text{-min}| < \delta \longrightarrow deriv\ (deriv\ f)\ y > 0)$ 
    by (metis assms(3,5) at-within-open cont-at-pos-imp-loc-pos' continuous-on-eq-continuous-within
      open-real-greaterThanLessThan)
  obtain  $\delta$  where  $\delta\text{-min}$ :  $\delta = \min\ \Delta\ (\min\ ((x\text{-min} - a) / 2)\ ((b - x\text{-min}) / 2))$ 
    by blast
  have  $\delta\text{-pos}$ :  $\delta > 0$ 
  proof (cases  $\delta = \Delta$ )
    show  $\delta = \Delta \implies 0 < \delta$ 
      by (simp add:  $\Delta\text{-pos}$ )
  next
    assume  $\delta \neq \Delta$ 
    then have  $\delta = \min\ ((x\text{-min} - a) / 2)\ ((b - x\text{-min}) / 2)$ 
      using  $\delta\text{-min}$  by linarith
    then show  $0 < \delta$ 
      using min-exists by force
  qed

```

```

have neigh-of-x-min-contained-in-ab:  $a < x - \text{min} - \delta \wedge x - \text{min} + \delta < b$ 
  by (smt (z3)  $\delta$ -min  $\delta$ -pos field-sum-of-halves)

have local-min:  $\forall x. |x - x - \text{min}| < \delta \longrightarrow f\ x \geq f\ x - \text{min}$ 
proof clarify
  fix x
  assume A:  $|x - x - \text{min}| < \delta$ 
  consider (eq)  $x = x - \text{min}$  | (lt)  $x < x - \text{min}$  | (gt)  $x > x - \text{min}$ 
    by linarith
  then show  $f\ x \geq f\ x - \text{min}$ 
  proof cases
    case eq
    then show ?thesis
      by simp
    next
    case lt
    have a-lt-x-and-xmin-lt-b:  $a < x \wedge x - \text{min} < b$ 
      using A neigh-of-x-min-contained-in-ab by linarith
    have  $f\ x > f\ x - \text{min}$ 
    proof (rule DERIV-neg-imp-decreasing-open[where  $a = x$ ])
      show  $x < x - \text{min}$ 
        by (simp add: lt)
    next
    fix y :: real
    assume x-lt-y:  $x < y$ 
    assume y-lt-x-min:  $y < x - \text{min}$ 
    — For  $x < x_{\min}$ , apply the Mean Value Theorem to  $f$  on  $[x, x_{\min}]$ .
    have  $\exists z > y. z < x - \text{min} \wedge \text{deriv}\ f\ x - \text{min} - \text{deriv}\ f\ y = (x - \text{min} - y) * \text{deriv}$ 
      (deriv f) z
    proof (rule MVT2[where  $a = y$  and  $b = x - \text{min}$  and  $f = \text{deriv}\ f$  and  $f' = \text{deriv}\ (\text{deriv}\ f)$ ])
      show  $y < x - \text{min}$ 
        by (simp add: y-lt-x-min)
    next
    fix z :: real
    assume y-lt-z:  $y \leq z$ 
    assume z-lt-x-min:  $z \leq x - \text{min}$ 
    show (deriv f has-real-derivative (deriv (deriv f)) z) (at z)
    proof (subst C2-cont-diff[where  $f = f$ , where  $U = \{a <..< b\}$ ])
      show C-k-on 2 f  $\{a <..< b\}$ 
        by (simp add: assms(2))
      show  $z \in \{a <..< b\}$  and True
        using a-lt-x-and-xmin-lt-b x-lt-y y-lt-z z-lt-x-min by auto
    qed
  qed
then obtain z where
  z-props:  $y < z < x - \text{min}$  and
  eq:  $\text{deriv}\ f\ x - \text{min} - \text{deriv}\ f\ y = (x - \text{min} - y) * \text{deriv}\ (\text{deriv}\ f)\ z$ 

```

```

    by blast
  have deriv f x-min = 0
    using fst-deriv-req by simp
  hence deriv f y = - (x-min - y) * deriv (deriv f) z
    using eq by linarith
  moreover have x-min - x > 0
    using lt by simp
  have deriv (deriv f) z > 0
    by (smt (verit) A Δ-prop δ-min x-lt-y z-props)
  ultimately have deriv f y < 0
    by (simp add: mult-less-0-iff y-lt-x-min)
  then show ∃ z. (f has-real-derivative z) (at y) ∧ z < 0
    by (meson C2-cont-diff a-lt-x-and-xmin-lt-b assms(2) dual-order.strict-trans
        greaterThanLessThan-iff x-lt-y y-lt-x-min)
next
  have continuous-on {a <..} f
    by (simp add: C2-cont-diff assms(2) differentiable-imp-continuous-on)
  then show continuous-on {x..x-min} f
    by (smt (verit, del-insts) a-lt-x-and-xmin-lt-b atLeastAtMost-iff
        continuous-on-subset greaterThanLessThan-iff subsetI)
qed
then show f x-min ≤ f x
  by simp
next
  case gt
  have a-lt-xmin-and-x-lt-b: a < x-min ∧ x < b
    using A ⟨a < x-min - δ ∧ x-min + δ < b⟩ by linarith
  have f x > f x-min
  proof (rule DERIV-pos-imp-increasing-open[where a = x-min])
    show x-min < x
      by (simp add: gt)
  next
    fix y :: real
    assume y-gt-xmin: x-min < y
    assume y-lt-x: y < x
    — For  $x_{\min} < y$ , apply the Mean Value Theorem to  $f'$  on  $[x_{\min}, y]$ .
    have ∃ z > x-min. z < y ∧ deriv f y - deriv f x-min = (y - x-min) * deriv
      (deriv f) z
    proof (rule MVT2[where a = x-min and b = y and f = deriv f and f'
      = deriv (deriv f)])
      show x-min < y
        by (simp add: y-gt-xmin)
    next
      fix z :: real
      assume z-ge-xmin: x-min ≤ z
      assume z-le-y: z ≤ y
      show (deriv f has-real-derivative (deriv (deriv f)) z) (at z)
      proof (subst C2-cont-diff[where f = f and U = {a<..}]
        show C-k-on 2 f {a<..})

```

```

      by (simp add: assms(2))
    show  $z \in \{a <..< b\}$  and True
      using a-lt-xmin-and-x-lt-b y-lt-x z-ge-xmin z-le-y by auto
    qed
  qed
  then obtain z where
    z-props:  $x\text{-min} < z < y$ 
    and eq:  $\text{deriv } f \ y - \text{deriv } f \ x\text{-min} = (y - x\text{-min}) * \text{deriv } (\text{deriv } f) \ z$ 
    by blast
  have  $\text{deriv } f \ x\text{-min} = 0$ 
    using fst-deriv-req by simp
  hence  $\text{deriv } f \ y = (y - x\text{-min}) * \text{deriv } (\text{deriv } f) \ z$ 
    using eq by simp
  moreover have  $y - x\text{-min} > 0$ 
    using y-gt-xmin by simp
  moreover have  $\text{deriv } (\text{deriv } f) \ z > 0$ 
    by (smt (verit, best) A  $\Delta$ -prop  $\delta$ -min y-lt-x z-props(1,2))
  ultimately have  $\text{deriv } f \ y > 0$ 
    by auto
  then show  $\exists d. (f \text{ has-real-derivative } d) \text{ (at } y) \wedge d > 0$ 
    by (meson C2-cont-diff a-lt-xmin-and-x-lt-b assms(2) dual-order.strict-trans
      greaterThanLessThan-iff y-lt-x y-gt-xmin)
next
  have continuous-on  $\{a <..< b\}$  f
    by (simp add: C2-cont-diff assms(2) differentiable-imp-continuous-on)
  then show continuous-on  $\{x\text{-min}..x\}$  f
    by (smt (verit, del-insts) a-lt-xmin-and-x-lt-b atLeastAtMost-iff
      continuous-on-subset greaterThanLessThan-iff subsetI)
  qed
  then show ?thesis
    by simp
  qed
  show ?thesis
    by (rule local-minimizer-from-neighborhood, smt  $\delta$ -pos local-min)
qed
end

```

5 Pathological Example: Non-Isolated Strict Local Minima

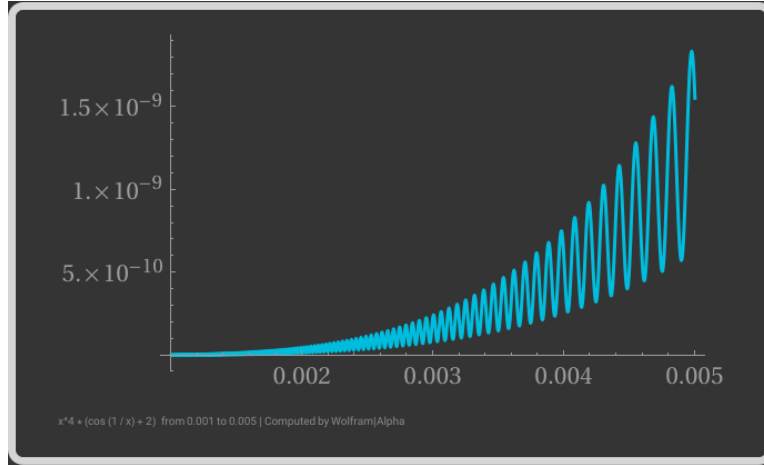
```

theory Cont-Nonisolated-Strict-Local-Minimizer-Exists
  imports Second-Derivative-Test HOL-Library.Quadratic-Discriminant
begin

```

Idea of the example. We construct a continuous function

$$f(x) = \begin{cases} x^4(\cos(1/x) + 2), & x \neq 0, \\ 0, & x = 0 \end{cases}$$



whose oscillations *speed up* as $x \rightarrow 0$ because of the $\cos(1/x)$ term. Multiplying by x^4 makes the function and its first derivative vanish at the origin, ensuring that $x = 0$ is a strict local minimizer, while the shifted cosine creates infinitely many additional strict local minimizers that accumulate at 0. Hence the minimizer at 0 is *strict* but *not isolated*.

theorem *Exists-Continuous-Func-with-non-isolated-strict-local-minimizer:*

$\exists f :: \text{real} \Rightarrow \text{real. continuous-on } \mathbb{R} \wedge$

$(\exists x\text{-star. strict-local-minimizer } f \ x\text{-star} \wedge \neg \text{isolated-local-minimizer } f \ x\text{-star})$

proof –

obtain f **where** $f\text{-def: } f = (\lambda(x :: \text{real}). \text{if } x \neq 0 \text{ then } x^4 * (\cos(1/x) + 2) \text{ else } 0)$

by *simp*

have $\text{deriv-}f: \wedge x :: \text{real. deriv } f \ x = (\text{if } x = 0 \text{ then } 0 \text{ else } x^2 * \sin(1/x) + 4 * x^3 * \cos(1/x) + 8 * x^3)$

$\wedge (\lambda x. f \ x) \text{ differentiable-on UNIV}$

$\wedge \text{deriv } (\text{deriv } f) \ x = (\text{if } x = 0 \text{ then } 0 \text{ else } 6*x * \sin(1/x) + (12*x^2 - 1)* \cos(1/x) + 24*x^2)$

$\wedge (\text{deriv } f) \text{ differentiable-on UNIV}$

proof (*safe*)

– First we compute the derivative away from 0, then we compute it at 0.

have $\text{deriv-}f\text{-at-nonzero:}$

$\wedge x. x \neq 0 \longrightarrow \text{deriv } f \ x = (x^2 * \sin(1/x) + 4*x^3 * \cos(1/x) + 8*x^3)$

$\wedge f \text{ field-differentiable at } x$

proof (*safe*)

fix $x :: \text{real}$

```

assume x-type: x ≠ 0

have cos-inverse-diff: (λw. cos (1 / w)) field-differentiable at x
proof -
  have f1: (λw. 1 / w) field-differentiable at x
    by (simp add: field-differentiable-divide x-type)
  have (λz. cos z) field-differentiable at (1 / x)
    by (simp add: field-differentiable-within-cos)
  then show ?thesis
    by (metis DERIV-chain2 f1 field-differentiable-def)
qed
then have (λx. cos (1 / x) + 2) field-differentiable at x
  by (simp add: Derivative.field-differentiable-add)
then have f2: (λx. x^4 * (cos (1 / x) + 2)) field-differentiable at x
  by (subst field-differentiable-mult, simp add: field-differentiable-power,
simp-all)

have deriv-2nd-part: deriv (λw. (λx. cos (1 / x) + 2) w) x = (sin (1 / x))
/ x^2
proof -
  have deriv (λw. (λx. cos (1 / x) + 2) w) x =
    (deriv (λw. (λx. cos (1 / x)) w) x + deriv (λw. (λx. 2) w) x)
  by (rule deriv-add, simp add: cos-inverse-diff, simp)
  also have ... = (sin (1 / x)) / x^2
  proof -
    have f1: DERIV (λz. cos z) (1 / x) :=> -sin (1 / x)
      by simp
    have f2: DERIV (λw. 1 / w) x :=> -1 / x^2
      using DERIV-inverse-func x-type by blast
    from f1 f2 have DERIV ((λz. cos z) ∘ (λw. 1 / w)) x :=> (-sin (1 / x))
    * (-1 / x^2)
      by (rule DERIV-chain)
    then show ?thesis
      by (simp add: DERIV-imp-deriv o-def)
  qed
  finally show ?thesis.
qed

show deriv f x = x^2 * sin (1 / x) + 4*x^3 * cos (1 / x) + 8*x^3
proof -
  have deriv f x = deriv (λx. x^4 * (cos (1 / x) + 2)) x
    by (metis (no-types, lifting) f-def mult-eq-0-iff power-zero-numeral)
  also have ... = x^4 * deriv (λx. cos (1 / x) + 2) x +
    deriv (λx. x^4) x * (cos (1 / x) + 2)
  by (rule deriv-mult, simp add: field-differentiable-power,
simp add: Derivative.field-differentiable-add cos-inverse-diff)
  also have ... = x^4 * (sin (1 / x)) / x^2 +
    deriv (λx. x^4) x * (cos (1 / x) + 2)
  by (simp add: deriv-2nd-part)

```

```

also have ... =  $x^4 * (\sin (1 / x)) / x^2 + (4*x^3) * (\cos (1 / x) + 2)$ 
  by (subst power-rule, simp)
also have ... =  $x^2 * (\sin (1 / x)) + (4*x^3) * (\cos (1 / x) + 2)$ 
  by (simp add: power2-eq-square power4-eq-xxxx)
also have ... =  $x^2 * \sin (1 / x) + 4*x^3 * \cos (1 / x) + 8*x^3$ 
  by (simp add: Rings.ring-distrib(2) mult.commute)
finally show ?thesis.
qed
from x-type f-def f2 show f field-differentiable at x
  by (subst field-differentiable-transfer-on-ball[where f =  $\lambda x. (x^4 * (\cos (1 / x) + 2))$ 
    / x) + 2))
    and  $\varepsilon = |x|$ , simp-all)
qed

have deriv-f-at-0: deriv f 0 = 0  $\wedge$  f field-differentiable at 0
proof -
  — By the definition of deriv, we need to show the limit of the difference quotient
  is 0.
  have dq-limit: (( $\lambda h. (f (0 + h) - f 0) / h$ )  $\longrightarrow$  0) (at 0)
  proof
    fix  $\varepsilon :: \text{real}$ 
    assume  $\varepsilon$ -pos:  $0 < \varepsilon$ 
    — Choose  $\delta > 0$  to make |difference quotient|  $< \varepsilon$ .
    obtain  $\delta$  where  $\delta$ -def:  $\delta = (\varepsilon / 3)$  powr (1 / 3)
      by simp
    — A reasonable  $\delta$  based on the growth of  $|h^3|$ .
    have  $\delta$ -pos:  $\delta > 0$ 
      using  $\varepsilon$ -pos by (simp add:  $\delta$ -def)
    have  $\exists \delta > 0. \forall h. 0 < |h| \wedge |h| < \delta \longrightarrow |(f (0 + h) - f 0) / h - 0| < \varepsilon$ 
    proof (intro exI[where  $x = \delta$ ], intro conjI insert  $\delta$ -pos, clarify)
      fix  $h :: \text{real}$ 
      assume  $h$ -pos:  $0 < |h|$ 
      assume  $h$ -lt- $\delta$ :  $|h| < \delta$ 

      have  $|(f (0 + h) - f 0) / h - 0| = |f h / h|$ 
        by (simp add: f-def)
      also have ... =  $|h^4 * (\cos (1 / h) + 2) / h|$ 
        using f-def by presburger
      also have ... =  $|h^3 * (\cos (1 / h) + 2)|$ 
      by (simp add: power3-eq-cube power4-eq-xxxx vector-space-over-itself.scale-scale)
      also have ...  $\leq |h^3| * |\cos (1 / h) + 2|$ 
        by (metis abs-mult order.refl)
      also have ...  $\leq |h^3| * (|\cos (1 / h)| + |2|)$ 
        by (simp add: mult-left-mono)
      also have ...  $\leq |h^3| * (1 + 2)$ 
        by (simp add: mult-left-mono)
      also have ... =  $3 * |h^3|$ 
        by simp
      also have ...  $< 3 * \delta^3$ 

```

```

    using power-strict-mono[of |h| δ 3] by (simp add: h-lt-δ power-abs)
  also have ... = 3 * (ε / 3)
    by (metis δ-def ε-pos div-self less-le more-arith-simps(5)
        mult-eq-0-iff pos-le-divide-eq powr-numeral powr-one-gt-zero-iff
        powr-powr times-divide-eq-left verit-comp-simplify(19)
        zero-neq-numeral)
    also have ... = ε
      by simp
    finally show |(f (0 + h) - f 0) / h - 0| < ε.
  qed
  then show ∃ d>0.∀ x∈UNIV. 0 < dist x 0 ∧ dist x 0 < d ⟶ dist ((f (0
+ x) - f 0) / x) 0 ≤ ε
    by (metis arithmetic-simps(57) dist-real-def less-le)
  qed
  then show ?thesis
    using DERIV-def DERIV-imp-deriv field-differentiable-def by blast
  qed

show deriv-f: ∧x. deriv f x =
  (if x = 0 then 0 else x2 * sin (1 / x) + 4 * x3 * cos (1 / x) + 8 * x3)
  using deriv-f-at-0 deriv-f-at-nonzero by presburger

show f-is-differentiable: (λx. f x) differentiable-on UNIV
  by (metis deriv-f-at-0 deriv-f-at-nonzero differentiable-on-def
      field-differentiable-imp-differentiable)

have snd-deriv-f-at-nonzero:
  ∧x. x ≠ 0 ⟶ deriv (deriv f) x = (6*x * sin (1 / x) + (12*x2 - 1)* cos
(1 / x) + 24*x2)
  ∧ (deriv f) field-differentiable at x
proof (safe)
  fix x :: real
  assume x-type: x ≠ 0

  have fst-term-diff: (λw. w2 * sin (1 / w)) field-differentiable at x
  proof -
    have f1: (λw. w2) field-differentiable at x
      by (simp add: field-differentiable-power)
    have (λw. sin (1 / w)) field-differentiable at x
      by (metis DERIV-chain2 DERIV-inverse-func field-differentiable-at-sin
          field-differentiable-def x-type)
    then show ?thesis
      by (simp add: f1 field-differentiable-mult)
  qed

  have fst-term-deriv: deriv (λw. w2 * sin (1 / w)) x = 2 * x * sin (1 / x)
  - cos (1 / x)
  proof -
    have deriv (λx. x2 * sin (1 / x)) x =

```

```

      x^2 * deriv (λx. sin (1 / x)) x + deriv (λx. x^2) x * sin (1 / x)
    by (rule deriv-mult, simp add: field-differentiable-power,
        metis DERIV-chain2 DERIV-inverse-func field-differentiable-at-sin
            field-differentiable-def x-type)
  moreover have deriv (λx. x^2) x = 2 * x
    using power-rule by auto
  moreover have deriv (λx. sin (1 / x)) x = -cos (1 / x) / x^2
  proof -
    have f1: DERIV (λz. sin z) (1 / x) :=> cos (1 / x)
      by simp
    have f2: DERIV (λx. 1 / x) x :=> -1 / x^2
      using DERIV-inverse-func x-type by blast
    from f1 f2 have DERIV ((λz. sin z) ∘ (λx. 1 / x)) x :=> cos (1 / x) *
(-1 / x^2)
      by (rule DERIV-chain)
    then show ?thesis
      by (simp add: DERIV-imp-deriv o-def)
  qed
  ultimately show ?thesis
    by (simp add: x-type)
  qed

  have snd-term-diff: (λx. 4 * x^3 * cos (1 / x)) field-differentiable at x
  proof -
    have t1: (λx. 4 * x^3) field-differentiable at x
      by (simp add: field-differentiable-power field-differentiable-mult)
    have t2: (λx. cos (1 / x)) field-differentiable at x
      by (metis DERIV-chain2 DERIV-inverse-func field-differentiable-at-cos
            field-differentiable-def x-type)
    show ?thesis
      by (simp add: t1 t2 field-differentiable-mult)
  qed
  have snd-term-diff': (λw. 4 * w^3 * cos (1 / w) + 8 * w^3) field-differentiable
at x
  proof -
    have t3: (λx. 8 * x^3) field-differentiable at x
      by (simp add: field-differentiable-mult field-differentiable-power)
    show ?thesis
      by (simp add: Derivative.field-differentiable-add t3 snd-term-diff)
  qed

  have snd-term-deriv:
    deriv (λx. 4 * x^3 * cos (1 / x) + 8 * x^3) x =
      12 * x^2 * cos (1 / x) + 4 * x * sin (1 / x) + 24 * x^2
  proof -
    have deriv (λx. 4 * x^3 * cos (1 / x) + 8 * x^3) x =
      deriv (λx. 4 * x^3 * cos (1 / x)) x + deriv (λx. 8 * x^3) x
    by (rule deriv-add, simp add: snd-term-diff,
        simp add: field-differentiable-mult field-differentiable-power)
  
```

also have ... = $(4 * x^3) * (\text{deriv } (\lambda x. \cos (1 / x)) x) +$
 $((12 * x^2) * (\cos (1 / x))) + \text{deriv } (\lambda x. 8 * x^3) x$

proof –

have $\text{deriv } (\lambda x. 4 * x^3 * \cos (1 / x)) x =$
 $(4 * x^3) * (\text{deriv } (\lambda x. \cos (1 / x)) x) +$
 $(\text{deriv } (\lambda x. 4 * x^3) x) * (\cos (1 / x))$

by (rule *deriv-mult*, *simp add: field-differentiable-mult field-differentiable-power*,
metis DERIV-fun-cos DERIV-inverse-func field-differentiable-def

x-type)

then have $\text{deriv } (\lambda x. 4 * x^3) x = 12 * x^2$

proof –

have $\text{deriv } (\lambda x. 4 * x^3) x = 4 * \text{deriv } (\lambda x. x^3) x$
by (rule *deriv-cmult*, *simp add: field-differentiable-power*)

then show ?thesis
by (*simp add: power-rule*)

qed

then show ?thesis

using $\langle \text{deriv } (\lambda x. 4 * x^3 * \cos (1 / x)) x = (4 * x^3) * (\text{deriv } (\lambda x. \cos$
 $(1 / x)) x) +$
 $(\text{deriv } (\lambda x. 4 * x^3) x) * (\cos (1 / x)) \rangle$

by *auto*

qed

also have ... = $(4 * x^3) * (\text{deriv } (\lambda x. \cos (1 / x)) x) +$
 $((12 * x^2) * (\cos (1 / x))) + 24 * x^2$

proof –

have $\text{deriv } (\lambda x. 8 * x^3) x = 24 * x^2$

proof –

have $\text{deriv } (\lambda x. 8 * x^3) x = 8 * \text{deriv } (\lambda x. x^3) x$
by (rule *deriv-cmult*, *simp add: field-differentiable-power*)

then show ?thesis
by (*simp add: power-rule*)

qed

then show ?thesis
by *auto*

qed

also have ... = $(4 * x^3) * \sin (1 / x) / x^2 + ((12 * x^2) * (\cos (1 / x)))$
 $+ 24 * x^2$

proof –

have $\text{deriv } (\lambda x. \cos (1 / x)) x = \sin (1 / x) / x^2$

proof –

have *f1*: $\text{DERIV } (\lambda z. \cos z) (1 / x) :=> -\sin (1 / x)$
by *simp*

have *f2*: $\text{DERIV } (\lambda x. 1 / x) x :=> -1 / x^2$
using *DERIV-inverse-func x-type* **by** *blast*

from *f1 f2* **have** $\text{DERIV } ((\lambda z. \cos z) \circ (\lambda x. 1 / x)) x :=> (-\sin (1 / x))$
 $* (-1 / x^2)$

by (rule *DERIV-chain*)

then show ?thesis
by (*simp add: DERIV-imp-deriv o-def*)

```

qed
then show ?thesis
  by auto
qed
also have ... = ((12 * x2) * (cos (1 / x))) + (4*x3) * sin (1 / x) / x2
+ 24 * x2
  by linarith
also have ... = (12 * x2) * (cos (1 / x)) + 4*x * sin (1 / x) + 24 * x2
proof -
  have (4*x3) * sin (1 / x) / x2 = 4*x * sin (1 / x)
    by (simp add: power2-eq-square power3-eq-cube)
  then show ?thesis
    by presburger
qed
finally show ?thesis.
qed

show deriv (deriv f) x = (6*x * sin (1 / x) + (12*x2 - 1)* cos (1 / x) +
24*x2)
proof -
  have deriv (deriv f) x = deriv (λx. x2 * sin (1 / x) + 4 * x3 * cos (1 /
x) + 8 * x3) x
  by (metis (no-types, opaque-lifting) deriv-f mult-cancel-left2 mult-cancel-right2
power-zero-numeral pth-7(2))
also have ... = deriv (λx. x2 * sin (1 / x) + (4 * x3 * cos (1 / x) + 8 *
x3)) x
  by (meson Groups.add-ac(1))
also have ... = deriv (λx. x2 * sin (1 / x)) x +
  deriv (λx. 4 * x3 * cos (1 / x) + 8 * x3) x
  by (rule deriv-add, simp add: fst-term-diff, simp add: snd-term-diff')
also have ... = 2 * x * sin (1 / x) - cos (1 / x) +
  deriv (λx. 4 * x3 * cos (1 / x) + 8 * x3) x
  by (simp add: fst-term-deriv)
also have ... = 2 * x * sin (1 / x) - cos (1 / x) +
  12 * x2 * cos (1 / x) + 4 * x * sin (1 / x) + 24 * x2
  by (simp add: snd-term-deriv)
also have ... = 2 * x * sin (1 / x) + 4 * x * sin (1 / x) +
  12 * x2 * cos (1 / x) - cos (1 / x) + 24 * x2
  by simp
also have ... = (6*x * sin (1 / x) + (12*x2 - 1)* cos (1 / x) + 24*x2)
  by (smt (verit, best) cos-add cos-zero mult-diff-mult sin-zero)
finally show ?thesis.
qed

show (deriv f) field-differentiable at x
proof (rule field-differentiable-transfer-on-ball
[where f = λ x. (x2 * sin (1 / x) + 4 * x3 * cos (1 / x) + 8 * x3)
and ε = |x|])

```

```

show 0 < |x|
  by (simp add: x-type)
show  $\forall y. y \in \text{ball } x \ |x| \longrightarrow y^2 * \sin (1 / y) + 4 * y^3 * \cos (1 / y) +$ 
 $8 * y^3 =$ 
  deriv f y
  by (simp add: deriv-f)
  show  $(\lambda x. x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3)$ field-differentiable at x
  by (simp add: Derivative.field-differentiable-add fst-term-diff is-num-normalize(1)

      snd-term-diff')
qed
qed

have deriv2-f-at-0:
  deriv (deriv f) 0 = 0  $\wedge$  (deriv f) field-differentiable at 0
proof -
  — By the definition of deriv, we need to show the limit of the difference
  quotient of  $f'$  is 0.
  have dq-limit:  $((\lambda h. (\text{deriv } f (0 + h) - \text{deriv } f 0) / h) \longrightarrow 0)$  (at 0)
  proof
    fix  $\varepsilon :: \text{real}$ 
    assume  $\varepsilon\text{-pos}$ : 0 <  $\varepsilon$ 
    have  $\exists \delta > 0. \forall h. 0 < |h| \wedge |h| < \delta \longrightarrow |(\text{deriv } f (0 + h) - \text{deriv } f 0) / h$ 
  — 0| <  $\varepsilon$ 
    proof (cases  $\varepsilon < 1/6$ )
      assume  $\text{eps-lt-inv6}$ :  $\varepsilon < 1/6$ 
      — Choose  $\delta > 0$  to ensure |difference quotient| <  $\varepsilon$ .
      obtain  $\delta$  where  $\delta\text{-def}$ :  $\delta = \varepsilon / 2$ 
      by blast
      have  $\delta\text{-pos}$ :  $\delta > 0$ 
      using  $\varepsilon\text{-pos}$  by (simp add:  $\delta\text{-def}$ )
      show  $\exists \delta > 0. \forall h. 0 < |h| \wedge |h| < \delta \longrightarrow |(\text{deriv } f (0 + h) - \text{deriv } f 0) /$ 
  h — 0| <  $\varepsilon$ 
    proof (intro exI[where  $x=\delta$ ], intro conjI insert  $\delta\text{-pos}$ , clarify)
      fix h :: real
      assume  $h\text{-pos}$ : 0 < |h|
      assume  $h\text{-lt-}\delta$ : |h| <  $\delta$ 

      have  $h\text{-bound1}$ : |h| <  $\varepsilon / 2$ 
      using  $h\text{-lt-}\delta$  by (simp add:  $\delta\text{-def}$ )
      have  $h\text{-bound2}$ :  $12 * |h^2| < \varepsilon / 2$ 
      proof -
        have |h| <  $\varepsilon / 2$  using  $h\text{-bound1}$  by blast
        then have  $|h^2| < (\varepsilon / 2)^2$ 
        by (metis abs-ge-zero abs-power2 power2-abs power-strict-mono
  zero-less-numeral)
        then have  $12 * |h^2| < 12 * (\varepsilon / 2)^2$ 
        by (simp add: mult-strict-left-mono)

```

```

also have ... = 12 * (ε2 / 4)
  by (simp add: power2-eq-square)
also have ... = 3 * ε2
  by simp
also have ... < ε/2
proof -
  have ε * 6 < 1
    by (meson eps-lt-inv6 less-divide-eq-numeral1(1))
  then show ?thesis
    by (simp add: ε-pos power2-eq-square)
qed
finally show ?thesis.
qed
have |(deriv f (0 + h) - deriv f 0) / h - 0| = |deriv f h / h|
  by (simp add: deriv-f-at-0)
also have ... = |(h2 * sin (1 / h) + 4*h3 * cos (1 / h) + 8*h3) / h|
  using deriv-f by presburger
also have ... = |(h2 * sin (1 / h) / h) + (4*h3 * cos (1 / h)) / h +
(8*h3) / h|
  by (simp add: add-divide-distrib)
also have ... = |h * sin (1 / h) + (4*h2 * cos (1 / h)) + 8 * h2|
  by (simp add: more-arith-simps(11) power2-eq-square power3-eq-cube)
also have ... ≤ |h * sin (1 / h)| + |4*h2 * cos (1 / h)| + |8 * h2|
  by linarith
also have ... ≤ |h| * |sin (1 / h)| + 4 * |h2| * |cos (1 / h)| + 8 * |h2|
  by (simp add: abs-mult)
also have ... ≤ |h| + 4 * |h2| + 8 * |h2|
proof -
  have i1: |h| * |sin (1 / h)| ≤ |h|
    using h-pos by fastforce
  have |h| * |cos (1 / h)| ≤ |h|
    by (simp add: mult-left-le)
  then show ?thesis
    by (smt (verit) cos-ge-minus-one cos-le-one i1 mult-left-le)
qed
also have ... = |h| + 12 * |h2|
  by simp
also have ... < ε
  using h-bound1 h-bound2 by auto
finally show |(deriv f (0 + h) - deriv f 0) / h - 0| < ε.
qed
next
assume ¬ ε < 1/6
then have ε ≥ 1/6 by linarith
then have eps-half: ε / 2 ≥ 1/12 by linarith
obtain δ where δ-def: δ = (1::real)/12 by blast
have δ-pos: δ > 0 using ε-pos by (simp add: δ-def)
show ∃ δ > 0. ∀ h. 0 < |h| ∧ |h| < δ → |(deriv f (0 + h) - deriv f 0) /
h - 0| < ε

```

```

proof (intro exI[where x=δ], intro conjI insert δ-pos, clarify)
  fix h :: real
  assume h-pos: 0 < |h|
  assume h-lt-δ: |h| < δ
  have h-bound1: |h| < ε / 2
  proof -
    have |h| < δ using h-lt-δ by blast
    also have ... = (1::real)/12 by (simp add: δ-def)
    also have ... ≤ ε / 2 using eps-half by blast
    finally show ?thesis.
  qed
  have h-bound2: 12 * |h|^2 < ε / 2
  proof -
    from h-bound1 have |h|^2 < (1/12)^2
    by (metis δ-def abs-ge-zero h-lt-δ power-strict-mono zero-less-numeral)
    hence 12 * |h|^2 < 12 * (1/12)^2
    by (rule mult-strict-left-mono, simp-all)
    also have ... = 1/12 by (simp add: power-one-over)
    also have ... ≤ ε / 2 using eps-half by blast
    finally show ?thesis.
  qed
  have |(deriv f (0 + h) - deriv f 0) / h - 0| = |deriv f h / h|
    by (simp add: deriv-f-at-0)
  also have ... = |(h^2 * sin (1 / h) + 4*h^3 * cos (1 / h) + 8*h^3) / h|
    using deriv-f by presburger
  also have ... = |(h^2 * sin (1 / h) / h) + (4*h^3 * cos (1 / h)) / h +
(8*h^3) / h|
    by (simp add: add-divide-distrib)
  also have ... = |h * sin (1 / h) + (4*h^2 * cos (1 / h)) + 8 * h^2|
    by (simp add: more-arith-simps(11) power2-eq-square power3-eq-cube)
  also have ... ≤ |h * sin (1 / h)| + |4*h^2 * cos (1 / h)| + |8 * h^2|
    by linarith
  also have ... ≤ |h| * |sin (1 / h)| + 4 * |h^2| * |cos (1 / h)| + 8 * |h^2|
    by (simp add: abs-mult)
  also have ... ≤ |h| + 4 * |h^2| + 8 * |h^2|
  proof -
    have i1: |h| * |sin (1 / h)| ≤ |h|
      using h-pos by fastforce
    have |h| * |cos (1 / h)| ≤ |h|
      by (simp add: mult-left-le)
    then show ?thesis
      by (smt (verit) cos-ge-minus-one cos-le-one i1 mult-left-le)
  qed
  also have ... = |h| + 12 * |h^2|
    by simp
  also have ... < ε
    using h-bound1 h-bound2 by auto
  finally show |(deriv f (0 + h) - deriv f 0) / h - 0| < ε.
qed

```

```

qed
then show  $\exists d > 0. \forall x \in UNIV. 0 < dist\ x\ 0 \wedge dist\ x\ 0 < d \longrightarrow$ 
       $dist\ ((deriv\ f\ (0 + x) - deriv\ f\ 0) / x)\ 0 \leq \varepsilon$ 
  by (metis cancel-comm-monoid-add-class.diff-zero dist-real-def le-less)
qed
then show ?thesis
  using DERIV-def DERIV-imp-deriv field-differentiable-def by blast
qed

show  $\bigwedge x. deriv\ (deriv\ f)\ x = (if\ x = 0\ then\ 0\ else\ 6 * x * sin\ (1 / x)$ 
       $+ (12 * x^2 - 1) * cos\ (1 / x)$ 
       $+ 24 * x^2)$ 
  using snd-deriv-f-at-nonzero deriv2-f-at-0 by presburger

show (deriv f) differentiable-on UNIV
  by (metis deriv2-f-at-0 differentiable-on-def
      field-differentiable-imp-differentiable snd-deriv-f-at-nonzero)
qed
then have f-cont: continuous-on  $\mathbb{R}$  f
  by (meson continuous-on-subset differentiable-imp-continuous-on top.extremum)
have f'-cont: continuous-on  $\mathbb{R}$  (deriv f)
  by (meson continuous-on-subset deriv-f differentiable-imp-continuous-on top.extremum)

obtain U where U-def:  $U = \{x :: real. -1 < x \wedge x < 1\}$ 
  by blast
then have open-neighborhood-of-zero:  $open\ U \wedge 0 \in U$ 
  using lemma-interval-lt by (subst open-dist, subst dist-real-def, fastforce)

have strict-local-minimizer-at-0: strict-local-minimizer f 0
  unfolding strict-local-minimizer-def strict-local-minimizer-on-def
proof (intro exI [where x=U], (subst sym[OF conj-assoc], rule conjI), rule open-neighborhood-of-zero)
  show  $\forall x \in U - \{0\}. f\ 0 < f\ x$ 
  proof
    fix x
    assume x-type:  $x \in U - \{0\}$ 
    then have x-nonzero:  $x \neq 0$ 
      by blast
    have  $cos(1/x) + 2 \geq 1$ 
      by (smt (verit) cos-ge-minus-one)
    then have  $x^4 * (cos(1/x) + 2) \geq x^4 * 1$ 
      by (rule mult-left-mono, force)
    then have  $f\ x \geq x^4$ 
      by (simp add: f-def x-nonzero)
    then have  $f\ x > 0$ 
      by (smt (verit, del-Insts) mult-le-0-iff power4-eq-xxxx x-nonzero zero-le-mult-iff)
    then show  $f\ 0 < f\ x$ 
      using f-def by force
  qed
qed
qed

```

then have *zero-min: local-minimizer f 0*
by (*simp add: strict-local-minimizer-imp-local-minimizer*)
have $(\exists x\text{-seq}::\text{nat} \Rightarrow \text{real}. (\forall n. \text{local-minimizer } f (x\text{-seq } n) \wedge x\text{-seq } n \neq 0) \wedge ((x\text{-seq} \longrightarrow 0) \text{ at-top}))$

proof –
obtain *left-seq :: nat \Rightarrow real where left-seq-def: $\forall n \in \mathbb{N}. n \neq 0 \longrightarrow$*
*left-seq n = inverse $((5 * \pi / 4) + 2 * \text{real } n * \pi)$*
by *force*
obtain *right-seq :: nat \Rightarrow real where right-seq-def: $\forall n \in \mathbb{N}. n \neq 0 \longrightarrow$*
*right-seq n = inverse $(\pi + 2 * \text{real } n * \pi)$*
by *force*

have *zero-lt-left-seq-lt-right-seq-both-pos: $\forall n \in \mathbb{N}. n \neq 0 \longrightarrow$*
 $0 < \text{left-seq } n \wedge \text{left-seq } n < \text{right-seq } n$

proof *clarify*
fix *n::nat*
assume *n-pos: $0 < n$*
then have *inv-left: inverse $(\text{left-seq } n) = (5 * \pi / 4) + 2 * \text{real } n * \pi$*
by (*metis bot-nat-0.not-eq-extremum id-apply inverse-inverse-eq left-seq-def of-nat-eq-id of-nat-in-Nats*)

have *inv-right: inverse $(\text{right-seq } n) = \pi + 2 * \text{real } n * \pi$*
by (*metis bot-nat-0.not-eq-extremum id-apply inverse-inverse-eq n-pos of-nat-eq-id of-nat-in-Nats right-seq-def*)

have *denom-ineq: $(\pi + 2 * \text{real } n * \pi) < ((5 * \pi / 4) + 2 * \text{real } n * \pi)$*
proof –
have $(5 * \pi / 4) + 2 * \text{real } n * \pi = 2 * \text{real } n * \pi + (5 * \pi / 4)$
by *simp*
have $((5 * \pi / 4) + 2 * \text{real } n * \pi) - (\pi + 2 * \text{real } n * \pi) =$
 $(5 * \pi / 4) + 2 * \text{real } n * \pi - \pi - 2 * \text{real } n * \pi$
by *simp*
also have $\dots = (5 * \pi / 4) - \pi$
by *simp*
also have $\dots = (5 * \pi / 4) - (4 * \pi / 4)$
by *simp*
also have $\dots = (5 - 4) * \pi / 4$
by *simp*
also have $\dots = \pi / 4$
by *simp*
then show *?thesis*
by *simp*

qed
then have *left-seq n < right-seq n*
by (*smt (verit) inv-left inv-right inverse-positive-iff-positive le-imp-inverse-le mult-nonneg-nonneg of-nat-less-0-iff pi-gt3*)

then show $0 < \text{left-seq } n \wedge \text{left-seq } n < \text{right-seq } n$
by (*smt (verit, best) denom-ineq inv-left inverse-positive-iff-positive mult-nonneg-nonneg*)

of-nat-less-0-iff pi-gt3)

qed

have *first-and-second-order-conditions*: $\forall n. n \neq 0 \longrightarrow$

$(\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. (y^2 * \sin(1/y) + 4 * y^3 * \cos(1/y) + 8 * y^3) = 0 \wedge$

$(6 * y * \sin(1/y) + (12 * y^2 - 1) * \cos(1/y) + 24 * y^2) > 0) \wedge$

$((\text{left-seq } n)^2 * \sin(1/(\text{left-seq } n)) + 4 * (\text{left-seq } n)^3 * \cos(1/(\text{left-seq } n)))$

$+ 8 * (\text{left-seq } n)^3 < 0 \wedge$

$((\text{right-seq } n)^2 * \sin(1/(\text{right-seq } n)) + 4 * (\text{right-seq } n)^3 * \cos(1/(\text{right-seq } n)))$

$+ 8 * (\text{right-seq } n)^3 > 0$

proof(*clarify*)

fix $n :: \text{nat}$

assume $n\text{-pos}$: $0 < n$

then have $n\text{-ge-1}$: $1 \leq n$

by *simp*

show $(\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. y^2 * \sin(1/y) + 4 * y^3 * \cos(1/y) + 8 * y^3 = 0 \wedge 0 < 6 * y * \sin(1/y) + (12 * y^2 - 1) * \cos(1/y) + 24 * y^2) \wedge$

$(\text{left-seq } n)^2 * \sin(1/\text{left-seq } n) + 4 * \text{left-seq } n^3 * \cos(1/\text{left-seq } n) + 8 * \text{left-seq } n^3 < 0 \wedge$

$0 < (\text{right-seq } n)^2 * \sin(1/\text{right-seq } n) + 4 * \text{right-seq } n^3 * \cos(1/\text{right-seq } n) + 8 * \text{right-seq } n^3$

proof *safe*

show *left-seq-less-zero*: $(\lambda x. x^2 * \sin(1/x) + 4 * x^3 * \cos(1/x) + 8 * x^3) (\text{left-seq } n) < 0$

proof $-$

obtain x **where** $x\text{-def}$: $x = \text{left-seq } n$

by *blast*

$-$ Rewrite $1/x$ in terms of $\frac{5\pi}{4} + 2n\pi$.

then have *inv-x-eqs*: $\text{inverse } x = \text{inverse} (\text{inverse} ((5 * \pi / 4) + 2 * \text{real } n * \pi))$

by (*metis bot-nat-0.not-eq-extremum id-apply left-seq-def n-pos of-nat-eq-id of-nat-in-Nats*)

then have $x\text{-inv}$: $1/x = (5 * \pi / 4) + 2 * \text{real } n * \pi$

by (*simp add: inverse-eq-divide*)

$-$ Evaluate $\sin(1/x)$ and $\cos(1/x)$.

have *sin-inv-x*: $\sin(1/x) = -(\text{sqrt } 2 / 2)$

proof $-$

have $\sin(1/x) = \sin((5 * \pi / 4) + 2 * \text{real } n * \pi)$

using $x\text{-inv}$ **by** *presburger*

also have $\dots = \sin(5 * \pi / 4)$

by (*simp add: sin-add*)

also have ... = - (sqrt 2 / 2)
using *sin-5pi-div-4* **by** *blast*
finally show $\sin (1 / x) = - (\text{sqrt } 2 / 2)$.
qed

have *cos-inv-x*: $\cos (1 / x) = - (\text{sqrt } 2 / 2)$
proof -
have *cos-val*: $\cos (1 / x) = \cos ((5 * \text{pi} / 4) + 2 * \text{real } n * \text{pi})$
using *x-inv* **by** *presburger*
also have ... = $\cos (5 * \text{pi} / 4)$
by (*simp add: cos-add*)
also have ... = - (sqrt 2 / 2)
using *cos-5pi-div-4* **by** *linarith*
finally show $\cos (1 / x) = - (\text{sqrt } 2 / 2)$.
qed

— Substitute these into the expression.

have *expr*: $x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3$
= - (sqrt 2 / 2) * x^2 + (8 - 2 * sqrt 2) * x^3

proof -
have $x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3$
= ($x^2 * - (\text{sqrt } 2 / 2)$) + $4 * x^3 * - (\text{sqrt } 2 / 2)$ + $8 * x^3$
by (*simp add: cos-inv-x sin-inv-x*)

also have ... = $x^2 * - (\text{sqrt } 2 / 2) + (-2 * \text{sqrt } 2) * x^3 + 8 * x^3$
by *simp*

also have ... = - (sqrt 2 / 2) * x^2 + (8 - 2 * sqrt 2) * x^3

proof -
have - (sqrt 2 / 2) + ($x^3 * (\text{sqrt } 2 * - 2) + x^3 * 8$) =
- (sqrt 2 / 2) + $x^3 * (\text{sqrt } 2 * - 2 + 8)$
by (*metis (no-types) nat-distrib(2)*)
then show *?thesis*
by (*simp add: Groups.mult-ac(2)*)

qed
finally show *rewrite-expr*:
 $x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3$
= - (sqrt 2 / 2) * x^2 + (8 - 2 * sqrt 2) * x^3 .

qed

— Factor out x^3 , and rewrite x^3 as $(\frac{5\pi}{4} + 2n\pi)^{-1}$.

have *deriv-right-seq-eval*: $\sin (1 / x) * x^2 + 4 * x^3 * \cos (1 / x) + 8$
* x^3 =
(- (sqrt 2 / 2) * ((5 * pi / 4) + 2 * real n * pi) + (8 - 2 * sqrt 2))
* x^3

proof -
have $\sin (1 / x) * x^2 + 4 * x^3 * \cos (1 / x) + 8 * x^3$ =
- (sqrt 2 / 2) * x^2 + (8 - 2 * sqrt 2) * x^3
by (*smt (verit, del-insts) Groups.mult-ac(2) cos-inv-x cos-zero di-*)

vide-eq-0-iff expr

*left-inverse more-arith-simps(11) one-power2 power2-eq-square
power3-eq-cube*

power-minus sin-inv-x sin-zero
also have ... = $(- (\text{sqrt } 2 / 2) * \text{inverse } x + (8 - 2 * \text{sqrt } 2)) * x^3$
by (*metis (no-types) distrib-right*)
also have ... = $(- (\text{sqrt } 2 / 2) * ((5 * \text{pi} / 4) + 2 * \text{real } n * \text{pi}) + (8 - 2 * \text{sqrt } 2)) * x^3$
by (*simp add: inv-x-eqs*)
finally show ?thesis.
qed

— Combine into a single fraction and show negativity.

have *first-term-eval*: $x^3 > 0$
by (*smt (verit) mult-nonneg-nonneg of-nat-0-le-iff pi-gt3 x-inv zero-compare-simps(7)*)

zero-less-power)
have *neg-term*: $(- (\text{sqrt } 2 / 2) * ((5 * \text{pi} / 4) + 2 * \text{real } n * \text{pi}) + (8 - 2 * \text{sqrt } 2)) < 0$
proof —
have *n-ge1*: $n \geq 1$
using *n-ge-1* **by** *auto*
have *lower-bound*: $2 * \text{real } n * \text{pi} \geq 2 * \text{pi}$
using *n-ge1* **by** (*simp add: mult-left-mono*)
then have *mult-bound*: $(- (\text{sqrt } 2 / 2) * ((5 * \text{pi} / 4) + 2 * \text{real } n * \text{pi}))$
 $\leq - (\text{sqrt } 2 / 2) * (5 * \text{pi} / 4 + 2 * \text{pi})$
by (*simp add: mult-left-mono*)
moreover have $(- (\text{sqrt } 2 / 2) * (5 * \text{pi} / 4 + 2 * \text{pi}) + (8 - 2 * \text{sqrt } 2)) < 0$
proof —
have $5 * \text{pi} / 4 + 2 * \text{pi} = 13 * \text{pi} / 4$
by *simp*
then have *simpification*: $(- (\text{sqrt } 2 / 2) * (5 * \text{pi} / 4 + 2 * \text{pi}) + (8 - 2 * \text{sqrt } 2))$
 $= (64 - 16 * \text{sqrt } 2 - 13 * \text{pi} * \text{sqrt } 2) / 8$
by (*simp add: field-simps*)
have *sufficies-to-show-numerator-neg*: $(64 - 16 * \text{sqrt } 2 - 13 * \text{pi} * \text{sqrt } 2) / 8 < 0$
 $= (64 - 16 * \text{sqrt } 2 - 13 * \text{pi} * \text{sqrt } 2 < 0)$
by *simp*
have $\text{sqrt } 2 * (16 + 13 * \text{pi}) > 64$
proof —
have *pi-gt-3*: $\text{pi} > 3$
by (*simp add: pi-gt3*)
hence $16 + 13 * \text{pi} > 16 + 13 * 3$
by (*simp add: mult-strict-left-mono*)
hence $16 + 13 * \text{pi} > 55$
by *simp*

```

then have  $\sqrt{2} * (16 + 13 * \pi) > \sqrt{2} * 55$ 
  by (simp add: mult-strict-left-mono)
moreover have  $\sqrt{2} * 55 > 64$ 
proof –
  have  $(\sqrt{2} * 55)^2 = 2 * 55^2$ 
    by (simp add: power-mult-distrib)
  also have  $\dots = 2 * (55 * 55)$ 
    by auto
  also have  $\dots = 6050$ 
    by simp
  also have  $\dots > 64 * 64$ 
    by eval
  moreover have  $\sqrt{2} * 55 > 0$ 
    by simp
  ultimately show  $\sqrt{2} * 55 > 64$ 
    using power-mono-iff
    by (metis less-le power2-eq-square zero-less-numeral)
qed
ultimately show ?thesis
  by linarith
qed
then have  $64 - 16 * \sqrt{2} - 13 * \pi * \sqrt{2} < 0$ 
  by (simp add: Groups.mult-ac(2) distrib-left)
then show ?thesis
  using simpification sufficies-to-show-numerator-neg by presburger
qed
then show ?thesis
  using mult-bound by linarith
qed
then show  $(\text{left-seq } n)^2 * \sin(1 / \text{left-seq } n) +$ 
 $4 * \text{left-seq } n^3 * \cos(1 / \text{left-seq } n) + 8 * \text{left-seq } n^3 < 0$ 
  by (metis deriv-right-seq-eval first-term-eval mult commute x-def
zero-compare-simps(10))
qed
show right-seq-greater-zero:  $(\lambda x. x^2 * \sin(1 / x) + 4 * x^3 * \cos(1 /$ 
 $x) + 8 * x^3)$ 
 $(\text{right-seq } n) > 0$ 
proof –
  obtain x where x-def:  $x = \text{right-seq } n$ 
    by blast
  then have inv-x-eqs:  $\text{inverse } x = \text{inverse}(\text{inverse}(\pi + 2 * \text{real } n * \pi))$ 
    by (metis id-apply n-pos of-nat-eq-id of-nat-in-Nats of-nat-less-0-iff
right-seq-def)
  have x-inv:  $1 / x = \pi + 2 * \text{real } n * \pi$ 
    unfolding right-seq-def by (metis inv-x-eqs inverse-eq-divide inverse-inverse-eq)

have sin-inv-x:  $\sin(1 / x) = 0$ 
  by (metis add.inverse-neutral sin-2npi sin-periodic-pi2 x-inv)

```

have *cos-inv-x*: $\cos (1 / x) = -1$
using *cos-2npi cos-periodic-pi2 x-inv* **by** *presburger*

have *f-x*: $x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3 = 4 * x^3$
by (*simp add: cos-inv-x sin-inv-x*)

have *x-pos*: $x > 0$
unfolding *right-seq-def*
by (*smt (verit) mult-nonneg-nonneg of-nat-less-0-iff pi-gt-zero x-inv zero-less-divide-iff*)

then show $0 < (\text{right-seq } n)^2 * \sin (1 / \text{right-seq } n) + 4 * \text{right-seq } n^3 * \cos (1 / \text{right-seq } n) + 8 * \text{right-seq } n^3$
using *cos-inv-x sin-inv-x x-def* **by** *fastforce*

qed

show $\exists y \in \{\text{left-seq } n.. \text{right-seq } n\}. y^2 * \sin (1 / y) + 4 * y^3 * \cos (1 / y) + 8 * y^3 =$
 $0 \wedge 0 < 6 * y * \sin (1 / y) + (12 * y^2 - 1) * \cos (1 / y) + 24 * y^2$

proof –

have *existence-of-minimizing-sequence*: $\exists y \in \{\text{left-seq } n.. \text{right-seq } n\}. y^2 * \sin (1 / y) + 4 * y^3 * \cos (1 / y) + 8 * y^3 = 0$

proof –

have $\exists x \geq \text{left-seq } n. x \leq \text{right-seq } n \wedge (\lambda x. x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3) x = 0$

proof(*rule IVT'*)

show $(\text{left-seq } n)^2 * \sin (1 / \text{left-seq } n) + 4 * \text{left-seq } n^3 * \cos (1 / \text{left-seq } n) + 8 * \text{left-seq } n^3 \leq 0$

using *left-seq-less-zero* **by** *auto*

show $0 \leq (\text{right-seq } n)^2 * \sin (1 / \text{right-seq } n) + 4 * \text{right-seq } n^3 * \cos (1 / \text{right-seq } n) + 8 * \text{right-seq } n^3$

using *right-seq-greater-zero* **by** *linarith*

show $\text{left-seq } n \leq \text{right-seq } n$

by (*metis id-apply leD linorder-linear n-pos of-nat-eq-id of-nat-in-Nats zero-lt-left-seq-lt-right-seq-both-pos*)

show *continuous-on* $\{\text{left-seq } n.. \text{right-seq } n\} (\lambda x. x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3)$

proof – — We prove continuity by establishing it is differentiable.

— First, note that left_seq_n is positive, so the interval does not contain 0.

have *left-seq-pos*: $\text{left-seq } n > 0$

by (*metis bot-nat-0.extremum-strict id-apply n-pos of-nat-eq-id of-nat-in-Nats zero-lt-left-seq-lt-right-seq-both-pos*)

— Transfer global differentiability to local differentiability of deriv *f*.

have $\bigwedge x. x \in \{\text{left-seq } n.. \text{right-seq } n\} \longrightarrow (\lambda x. x^2 * \sin (1 / x) + 4 * x^3 * \cos (1 / x) + 8 * x^3) \text{ field-differentiable at } x$

```

proof clarify
  fix x::real
  assume x-type: x ∈ {left-seq n..right-seq n}
  show (λx. x2 * sin (1 / x) + 4 * x3 * cos (1 / x) + 8 * x3)
field-differentiable at x
  proof(rule field-differentiable-transfer-on-ball[where f = deriv f
and ε = x])
  show 0 < x
  using left-seq-pos x-type by auto
  show ∀ y. y ∈ ball x x → deriv f y = y2 * sin (1 / y) + 4 * y3 *
cos (1 / y) + 8 * y3
  by (simp add: deriv-f)
  show deriv f field-differentiable at x
  by (meson UNIV-I deriv-f differentiable-on-def field-differentiable-def
real-differentiableE)
  qed
qed
  then have (λx. x2 * sin (1 / x) + 4 * x3 * cos (1 / x) + 8 * x3)
differentiable-on {left-seq n..right-seq n}
  by (meson differentiable-at-imp-differentiable-on field-differentiable-imp-differentiable)
  then show ?thesis
  using differentiable-imp-continuous-on by blast
  qed
qed
  then show ∃ y∈{left-seq n..right-seq n}. y2 * sin (1 / y) + 4 * y3 *
cos (1 / y) + 8 * y3 = 0
  by presburger
  qed
  then obtain min-n where min-n-def: min-n ∈ {left-seq n..right-seq n} ∧
min-n2 * sin (1 / min-n) + 4 * min-n3 * cos (1 / min-n) + 8 * min-n3 =
0
  by blast
  have ∧ y. y ∈ {left-seq n .. right-seq n} → 0 < 6 * y * sin (1 / y) +
(12 * y2 - 1) * cos (1 / y) + 24 * y2
  proof (clarify)
  fix y :: real
  assume y-int: y ∈ {left-seq n .. right-seq n}
  — Since left_seq_n > 0, every y in the interval is positive.
  then have y-pos: y > 0
  by (metis atLeastAtMost-iff bot-nat-0.extremum id-apply linorder-not-less
n-pos
of-nat-eq-id of-nat-in-Nats order-less-le-trans zero-lt-left-seq-lt-right-seq-both-pos)

  have ∃ x-nc :: real ⇒ real. ∀ c ∈ {0..pi/4}. x-nc c = inverse (pi + c
+ 2*pi*real n)
  by auto
  then obtain x-nc :: real ⇒ real where x-nc-def: ∀ c ∈ {0..pi/4}. x-nc
c = inverse (pi + c + 2*pi*real n)
  by auto

```

```

have  $\exists x\text{-nc} :: \text{real} \Rightarrow \text{real}. \forall c \in \{0..pi/4\}. x\text{-nc } c = \text{inverse } (pi + c$ 
 $+ 2*pi*real\ n)$ 
by auto
then obtain  $x\text{-nc} :: \text{real} \Rightarrow \text{real}$  where  $x\text{-nc-def}: \forall c \in \{0..pi/4\}. x\text{-nc}$ 
 $c = \text{inverse } (pi + c + 2*pi*real\ n)$ 
by auto
have continuous-on-subinterval: continuous-on  $\{0..pi/4\}$   $x\text{-nc}$ 
proof –
have cont-denom: continuous-on  $\{0..pi/4\}$   $(\lambda c. pi + c + 2*pi*real\ n)$ 
proof –
have continuous-on  $\{0..pi/4\}$   $(\lambda c. c)$ 
using continuous-on-id by blast
moreover have continuous-on  $\{0..pi/4\}$   $(\lambda c. pi + 2*pi*real\ n)$ 
using continuous-on-const by blast
ultimately show ?thesis
by (simp add: continuous-on-add)
qed
then have continuous-on  $\{0..pi/4\}$   $(\lambda x. \text{inverse } ((\lambda c. pi + c +$ 
 $2*pi*real\ n)\ x))$ 
by(rule continuous-on-inverse,
smt (verit) add-mono-thms-linordered-field(4) atLeastAtMost-iff
of-nat-less-0-iff pi-neq-zero pi-not-less-zero zero-compare-simps(4)))

then show ?thesis
using continuous-on-cong x-nc-def by fastforce
qed

have minimizer-dom:  $\exists x. 0 \leq x \wedge x \leq pi/4 \wedge x\text{-nc } x = y$ 
proof(rule IVT2')
show  $x\text{-nc } (pi / 4) \leq y$ 
proof –
have  $x\text{-nc } (pi / 4) = \text{inverse } (pi + pi / 4 + 2 * real\ n * pi)$ 
by (metis (no-types, opaque-lifting) atLeastAtMost-iff divide-eq-imp
divide-real-def linorder-not-less mult.left-commute mult.right-neutral

mult-le-0-iff nle-le of-nat-0-le-iff of-nat-numeral pi-gt-zero x-nc-def

zero-neq-numeral)
also have  $\dots = \text{inverse } ((5 * pi / 4) + 2 * real\ n * pi)$ 
by simp
also have  $\dots = \text{left-seq } n$ 
by (metis bot-nat-0.not-eq-extremum id-apply left-seq-def n-pos
of-nat-eq-id of-nat-in-Nats)
also have  $\dots \leq y$ 
using y-int by presburger
finally show ?thesis.
qed
show  $y \leq x\text{-nc } 0$ 
proof –

```

```

have  $y \leq \text{right\_seq } n$ 
  using  $y\text{-int}$  by  $\text{presburger}$ 
also have  $\dots = \text{inverse } (\pi + 2 * \text{real } n * \pi)$ 
  by ( $\text{metis bot-nat-0.not-eq-extremum id-apply n-pos of-nat-eq-id}$ 
 $\text{of-nat-in-Nats right-seq-def}$ )
also have  $\dots = x\text{-nc } 0$ 
  using  $x\text{-nc-def}$  by  $\text{auto}$ 
finally show  $?thesis$ .
qed
show  $0 \leq \pi / 4$ 
  by  $\text{simp}$ 
show  $\text{continuous-on } \{0.. \pi / 4\} x\text{-nc}$ 
  using  $\text{continuous-on-subinterval}$  by  $\text{simp}$ 
qed
then have  $\text{minimizer-dom}' : \exists c \in \{0.. \pi / 4\}. y = x\text{-nc } c$ 
  using  $\text{atLeastAtMost-iff}$  by  $\text{blast}$ 

```

— We will show that $f''(x_{nc}(c)) > 0$ for all $c \in [0, 1]$, then use the fact that $\text{left_seq}_n \leq x_{nc}(c) \leq \text{right_seq}_n$ together with the IVT to establish the existence of $c \in [0, \frac{\pi}{4}]$ such that $x_{nc}(c) = y$, and then conclude that $f''(y) > 0$.

```

have  $\text{snd-deriv-positive-in-neighborhood} : \forall c \in \{0.. \pi / 4\}. \text{left-seq } n \leq$ 
 $x\text{-nc } c \wedge x\text{-nc } c \leq \text{right-seq } n \wedge \text{deriv } (\text{deriv } f) (x\text{-nc } c) > 0$ 
proof ( $\text{safe}$ )
  fix  $c :: \text{real}$ 
  assume  $c\text{-type} : c \in \{0.. \pi / 4\}$ 
  then have  $c\text{-bounds} : 0 \leq c \wedge c \leq \pi / 4$ 
    by  $\text{simp}$ 

  have  $x\text{-nc-eqs} : x\text{-nc } c = \text{inverse } (\pi + c + 2 * \pi * \text{real } n)$ 
    using  $c\text{-bounds}$   $\text{inverse-eq-divide}$   $\pi\text{-half-le-two}$   $x\text{-nc-def}$  by  $\text{auto}$ 
  show  $\text{left-seq } n \leq x\text{-nc } c$ 
proof —
  have  $f1 : \text{left-seq } n = \text{inverse } ((5 * \pi / 4) + 2 * \text{real } n * \pi)$ 
    by ( $\text{metis bot-nat-0.not-eq-extremum id-apply left-seq-def n-pos}$ 
 $\text{of-nat-eq-id of-nat-in-Nats}$ )
  from  $c\text{-bounds}$  have  $1 / ((5 * \pi / 4) + 2 * \text{real } n * \pi) \leq 1 / (\pi +$ 
 $c + 2 * \pi * \text{real } n)$ 
    by ( $\text{subst frac-le, simp-all, simp add: add-sign-intros(1)}$ )
  then show  $?thesis$ 
    by ( $\text{simp add: f1 x-nc-eqs inverse-eq-divide}$ )
qed

then have  $x\text{-nc-pos} : x\text{-nc } c > 0$ 
  by ( $\text{metis id-apply n-pos of-nat-eq-id of-nat-in-Nats order-less-le-trans}$ 
 $\text{zero-lt-left-seq-lt-right-seq-both-pos zero-order(5)}$ )

show  $x\text{-nc } c \leq \text{right-seq } n$ 
proof —

```

```

      have f1: right-seq n = inverse (pi + 2 * real n * pi)
      by (metis bot-nat-0.not-eq-extremum id-apply n-pos of-nat-eq-id
of-nat-in-Nats right-seq-def)
      from c-bounds have 1 / (pi + c + 2*pi*real n) ≤ 1 / (pi + 2 * real
n * pi)
      by (subst frac-le, simp-all, smt (verit, del-insts) m2pi-less-pi
mult-sign-intros(1) of-nat-less-0-iff)
      then show ?thesis
      by (simp add: f1 x-nc-eqs inverse-eq-divide)
    qed

— Bounds on sin(c) and cos(c).
have pi + c + 2*pi*real n ≥ 3*pi
proof —
  have pi + c + 2*pi*real n ≥ pi + 0 + 2*pi*real 1
  by (smt (verit, best) Num.of-nat-simps(2) c-bounds mult-left-mono
n-ge-1
      pi-not-less-zero real-of-nat-ge-one-iff)
  then show ?thesis
  by linarith
  qed
then have x-nc-bound: x-nc c ≤ inverse(3*pi)
  by (smt (verit) le-imp-inverse-le pi-gt-zero x-nc-eqs)
  then have cos-coef-bound: (1 - 12 * (x-nc c)^2) ≥ (1 - 12 *
(inverse(3*pi))^2)
  using x-nc-pos by force

have sin-bound: 0 ≤ sin c ∧ sin c ≤ sqrt(2)/2
proof safe
  show 0 ≤ sin c
  using c-bounds sin-ge-zero by auto
  show sin c ≤ sqrt(2)/2
  by (smt (verit, best) c-bounds frac-le pi-not-less-zero sin-45
sin-mono-less-eq)
  qed
have cos-bound: sqrt(2)/2 ≤ cos c ∧ cos c ≤ 1
proof safe
  show sqrt 2 / 2 ≤ cos c
  by (smt (verit) c-bounds cos-45 cos-monotone-0-pi-le machin
pi-machin)
  show cos c ≤ 1
  by simp
  qed

show 0 < deriv (deriv f) (x-nc c)
proof —
  — Lower bound of f''(xnc).
  have snd-deriv-at-x-nc: deriv (deriv f) (x-nc c) = (1 - 12 * (x-nc
c)^2) * cos c - 6 * (x-nc c) * sin c + 24 * (x-nc c)^2

```

```

proof–
  have f1:  $\sin (1 / (x-nc\ c)) = -\sin\ c$ 
  proof –
    have  $\sin (1 / (x-nc\ c)) = \sin (pi + c + 2*pi*real\ n)$ 
    by (simp add: inverse-eq-divide x-nc-eqs)
    also have ... =  $\sin (pi + c)$ 
    by (metis Groups.mult-ac(2) id-apply-of-real-eq-id sin.plus-of-nat)
    also have ... =  $-\sin\ c$ 
    by simp
    finally show ?thesis.
  qed
  have f2:  $\cos (1 / (x-nc\ c)) = -\cos\ c$ 
  proof –
    have  $\cos (1 / (x-nc\ c)) = \cos (pi + c + 2*pi*real\ n)$ 
    by (simp add: inverse-eq-divide x-nc-eqs)
    also have ... =  $\cos (pi + c)$ 
    by (metis Groups.mult-ac(2) id-apply-of-real-eq-id cos.plus-of-nat)
    also have ... =  $-\cos\ c$ 
    by simp
    finally show ?thesis.
  qed

  have  $deriv\ (deriv\ f)\ (x-nc\ c) = (12*(x-nc\ c)^2 - 1)*\ \cos (1 / (x-nc\ c)) + 6*(x-nc\ c) * \sin (1 / (x-nc\ c)) + 24*(x-nc\ c)^2$ 
    using deriv-f x-nc-pos by auto
  also have ... =  $(1 - 12 * (x-nc\ c)^2) * \cos\ c - 6 * (x-nc\ c) * \sin\ c + 24 * (x-nc\ c)^2$ 
    by (smt (verit) f1 f2 minus-mult-commute more-arith-simps(8))
  finally show ?thesis.
  qed
  have snd-deriv-bound:  $deriv\ (deriv\ f)\ (x-nc\ c) \geq (1 - 12 * (x-nc\ c)^2 - 6 * (x-nc\ c)) * (\sqrt{2} / 2)$ 
  proof –
    have  $deriv\ (deriv\ f)\ (x-nc\ c) \geq (1 - 12 * (x-nc\ c)^2) * \cos\ c - 6 * (x-nc\ c) * (\sqrt{2}/2) + 24 * (x-nc\ c)^2$ 
    using snd-deriv-at-x-nc sin-bound x-nc-pos by auto
    also have ...  $\geq (1 - 12 * (x-nc\ c)^2 - 6 * (x-nc\ c)) * (\sqrt{2} / 2)$ 
    by (smt (verit, best) calculation cos-bound divide-pos-pos one-power2 real-le-rsqrt right-diff-distrib' sum-le-prod1 vector-space-over-itself.scale-left-diff-distrib zero-compare-simps(12))
    then show ?thesis.
  qed

  show  $0 < deriv\ (deriv\ f)\ (x-nc\ c)$ 
  proof –
    obtain h :: real  $\Rightarrow$  real where h-def:  $h = (\lambda x. - 12 * x^2 - 6 * x + 1)$ 
    by auto
    have diff-h:  $\forall x. h$  field-differentiable at x

```

```

      unfolding h-def
    proof clarify
      fix x::real
      have d1: ( $\lambda x. -12 * x^2$ ) field-differentiable at x
    by(rule field-differentiable-mult, simp, simp add: field-differentiable-power)
      have d2: ( $\lambda x. -6 * x$ ) field-differentiable at x
    by(rule field-differentiable-mult, simp, simp add: field-differentiable-power)
      from d1 d2 show ( $\lambda x. -12 * x^2 - 6 * x + 1$ ) field-differentiable
at x
    by(subst field-differentiable-add, simp add: Derivative.field-differentiable-diff,
simp-all)
      qed

      have h-roots:  $\forall x. h x = 0 \iff x = (-6 + \text{sqrt } 84) / 24 \vee x =$ 
 $(-6 - \text{sqrt } 84) / 24$ 
      proof(clarify)
        fix x ::real
        have roots:  $(12 * x^2 + 6 * x + (-1) = 0) = (x = (-6 + \text{sqrt}$ 
 $(6^2 - 4 * 12 * (-1))) / (2 * 12) \vee x = (-6 - \text{sqrt}(6^2 - 4 * 12 * (-1))) /$ 
 $(2 * 12))$ 
          using discrim-def by(subst discriminant-iff, eval, force)

        then show  $(h x = 0) = (x = (-6 + \text{sqrt } 84) / 24 \vee x = (-6$ 
 $- \text{sqrt } 84) / 24)$ 
          using h-def by auto
        qed

        have right-root-positive:  $(-6 + \text{sqrt } 84) / 24 > 0$ 
        proof -
          have  $-6 + \text{sqrt } 84 > -6 + \text{sqrt } 64$ 
            by (smt (verit) real-sqrt-less-mono)
          then show  $(-6 + \text{sqrt } 84) / 24 > 0$ 
            by simp
          qed

        then have left-root-neg:  $0 > (-6 - \text{sqrt } 84) / 24$ 
          by fastforce
        have h-pos-on-interval:  $\forall x \in \{0..<(-6 + \text{sqrt } 84) / 24\}. h x > 0$ 
        proof(rule ccontr)
          assume  $\neg (\forall x \in \{0..<(-6 + \text{sqrt } 84) / 24\}. 0 < h x)$ 
          then obtain z where z-def:  $z \in \{0..<(-6 + \text{sqrt } 84) / 24\} \wedge$ 
 $0 \geq h z$ 
            by fastforce
          then have z-not-root:  $z \neq (-6 + \text{sqrt } 84) / 24 \wedge z \neq (-6 -$ 
 $\text{sqrt } 84) / 24$ 
            using z-def by force
          show False
        proof(cases h z = 0)
          show  $h z = 0 \implies \text{False}$ 
            using h-roots z-not-root by blast
        qed
      end

```

```

next
  assume  $h\ z \neq 0$ 
  then have  $hz\text{-neg}: h\ z < 0$ 
    using  $z\text{-def}$  by auto
  have  $\exists x. 0 \leq x \wedge x \leq z \wedge h\ x = 0$ 
  proof(rule IVT2')
    show  $h\ z \leq 0$ 
      by (simp add:  $z\text{-def}$ )
    show  $0 \leq h\ 0$ 
      by (simp add:  $h\text{-def}$ )
    show  $0 \leq z$ 
      using  $z\text{-def}$  by fastforce
    show continuous-on  $\{0..z\}$   $h$ 
      by (meson continuous-at-imp-continuous-on diff-h
field-differentiable-imp-continuous-at)
  qed
  then show False
  by (metis atLeastLessThan-iff h-roots left-root-neg not-less  $z\text{-def}$ )
  qed
qed

have  $(-6 + \sqrt{84}) / 24 > 1 / (3 * \pi)$ 
proof -
  have  $i1: 64 / \pi^2 < 8$ 
  proof -
    have  $\pi * \pi > 3 * 3$ 
  by (meson  $\pi\text{-gt}3$  mult-strict-mono  $\pi\text{-gt-zero}$  verit-comp-simplify( $\gamma$ ))
  then have  $\pi^2 > 9$ 
    by (simp add: power2-eq-square)
  then have  $64 / \pi^2 < 64 / 8$ 
    by (smt (verit) frac-less2)
  also have  $\dots = 8$ 
    by eval
  finally show ?thesis.
  qed

have  $i2: 96 / \pi < 32$ 
proof -
  have  $96 / \pi < 96 / 3$ 
    by (meson frac-less2 order.refl  $\pi\text{-gt}3$  verit-comp-simplify(19))
  also have  $\dots = 32$ 
    by eval
  finally show ?thesis.
  qed

have  $(8 / \pi + 6)^2 < 84$ 
proof -
  have  $((8::\text{real}) / \pi + 6)^2 = (8 / \pi)^2 + 2 * (8 / \pi) * 6 + 6^2$ 
    by (simp add: power2-sum)

```

```

    also have ... = 82/pi2 + 2*(8/pi)*6 + 62
      by (simp add: power-divide)
    also have ... = 64/pi2 + 96/pi + 36
      by simp
    also have ... < 84
      using i1 i2 by linarith
    finally show ?thesis.
  qed
  then have lt-sqrt84: 8/pi + 6 < sqrt(84)
    using real-less-rsqrt by presburger
  have lt-3pi-sqrt84: 24 + 18*pi < 3 * pi * sqrt(84)
  proof -
    have 24 + 18*pi = 3*8 + 3*6*pi
      by simp
    also have ... = 3*pi*(8/pi) + 3*pi*6
      by simp
    also have ... = 3*pi*((8/pi)+6)
      by (simp add: distrib-left)
    also have ... < 3 * pi * sqrt(84)
      by (simp add: lt-sqrt84)
    finally show ?thesis.
  qed
  have (-6+sqrt(84))*(3*pi) > 24
  proof -
    have (-6+sqrt(84))*(3*pi) = -6*(3*pi) + sqrt(84)*(3*pi)
      by (meson ring-class.ring-distrib(2))
    also have ... = -18*pi + 3*pi * sqrt(84)
      by simp
    also have ... > 24
      using lt-3pi-sqrt84 by auto
    finally show ?thesis.
  qed
  then have (-6+sqrt(84))*(3*pi) / 24 > 1
    by simp
  then show (-6+sqrt(84)) / 24 > 1 / (3*pi)
    by (metis pi-gt-zero pos-divide-less-eq times-divide-eq-left
    zero-compare-simps(6) zero-less-numeral)
  qed
  then have x-nc c < (-6+sqrt(84)) / 24
    by (metis dual-order.strict-trans2 inverse-eq-divide x-nc-bound)
  then have h-x-nc-pos: h (x-nc c) > 0
    by (simp add: h-pos-on-interval less-eq-real-def x-nc-pos)

  have deriv (deriv f) (x-nc c) ≥ (sqrt(2)/2) * h (x-nc c)
    by (metis Groups.mult-ac(2) snd-deriv-bound diff-add-eq h-def
    mult-minus-left uminus-add-conv-diff)
  then show ?thesis
  by (smt (verit) h-x-nc-pos half-gt-zero-iff mult-pos-pos real-sqrt-gt-0-iff)
  qed

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      qed
    qed

    then show  $0 < 6 * y * \sin (1 / y) + (12 * y^2 - 1) * \cos (1 / y) + 24 * y^2$ 
      by (smt (verit, best) deriv-f minimizer-dom')
    qed
    then show  $\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. y^2 * \sin (1 / y) + 4 * y^3 * \cos (1 / y) + 8 * y^3 = 0 \wedge 0 < 6 * y * \sin (1 / y) + (12 * y^2 - 1) * \cos (1 / y) + 24 * y^2$ 
      using min-n-def by blast
    qed
  qed
  qed
  qed

  have optimality-conditions:  $\forall n. n \neq 0 \longrightarrow (\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. (\text{deriv } f) y = 0 \wedge \text{deriv } (\text{deriv } f) y > 0)$ 
  proof clarify
    fix n::nat
    assume  $0 < n$ 
    then obtain min-n where min-n-def:  $\text{min-n} \in \{\text{left-seq } n .. \text{right-seq } n\} \wedge \text{min-n}^2 * \sin (1 / \text{min-n}) + 4 * \text{min-n}^3 * \cos (1 / \text{min-n}) + 8 * \text{min-n}^3 = 0 \wedge 0 < 6 * \text{min-n} * \sin (1 / \text{min-n}) + (12 * \text{min-n}^2 - 1) * \cos (1 / \text{min-n}) + 24 * \text{min-n}^2$ 
      using first-and-second-order-conditions bot-nat-0.not-eq-extremum by presburger
    have fst-order-condition:  $\text{deriv } f \text{ min-n} = 0$ 
      using deriv-f min-n-def by presburger
    have snd-order-condition:  $\text{deriv } (\text{deriv } f) \text{ min-n} > 0$ 
      using deriv-f min-n-def by fastforce
    show  $\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. \text{deriv } f y = 0 \wedge 0 < \text{deriv } (\text{deriv } f) y$ 
      using fst-order-condition min-n-def snd-order-condition by blast
    qed

  have seq-of-local-minizers-exists:  $\forall n. n \neq 0 \longrightarrow (\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. \text{local-minimizer } f y)$ 
  proof (clarify)
    fix n::nat
    assume n-pos:  $0 < n$ 
    then obtain y where y-def:  $(y \in \{\text{left-seq } n .. \text{right-seq } n\} \wedge (\text{deriv } f) y = 0 \wedge \text{deriv } (\text{deriv } f) y > 0)$ 
      using gr-implies-not0 optimality-conditions by presburger
    have right-seq-def2:  $\text{right-seq } n = \text{inverse } (\pi + 2 * \text{real } n * \pi)$ 
      by (metis id-apply less-not-refl n-pos of-nat-eq-id of-nat-in-Nats right-seq-def)

    have  $y \in \{\text{left-seq } n .. \text{right-seq } n\} \wedge \text{local-minimizer } f y$ 
    proof (subst second-derivative-test[where a = left-seq n, where b = right-seq n])

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show proper-interval: left-seq n < right-seq n
by (metis (no-types) id-apply n-pos of-nat-eq-id of-nat-in-Nats rel-simps(70)
zero-lt-left-seq-lt-right-seq-both-pos)
show C-k-on 2 f {left-seq n <..proof(rule C-k-on-subset[where U = {0<..<(1::real)}])
  show f-contin-diff-on-right: C-k-on 2 f {0<..<(1::real)}
  proof(rule C2-on-open-U-def2)
    show open {0<..<(1::real)}
    using lemma-interval by(subst open-dist, subst dist-real-def, simp add:
abs-minus-commute lemma-interval-lt)
    show f differentiable-on {0<..<(1::real)}
      by (meson deriv-f differentiable-on-subset top.extremum)
    show deriv f differentiable-on {0<..<(1::real)}
      by (meson deriv-f differentiable-on-subset top.extremum)
    show continuous-on {0<..<(1::real)} (deriv (deriv f))
    proof -
      have  $\forall x \in \{0 < .. < 1\}. \text{deriv (deriv f) } x = 6*x * \sin(1/x) + (12*x^2 - 1)*\cos(1/x) + 24*x^2$ 
      by (simp add: deriv-f)
      moreover have continuous-on {0<..<(1::real)} ( $\lambda x. 6*x * \sin(1/x) + (12*x^2 - 1)*\cos(1/x) + 24*x^2$ )
      proof -
        have  $\{0 < .. < (1::\text{real})\} \subseteq \{x :: \text{real}. x > 0\}$ 
        by fastforce
        moreover have continuous-on  $\{x :: \text{real}. x > 0\}$  ( $\lambda x. 6*x * \sin(1/x) + (12*x^2 - 1)*\cos(1/x) + 24*x^2$ )
        by (auto intro!: continuous-intros)
        ultimately show ?thesis
        using continuous-on-subset by blast
      qed
    ultimately show continuous-on {0<..<1} (deriv (deriv f))
    using continuous-on-cong by fastforce
  qed
qed

show open {left-seq n <..\wedge {left-seq n <..\subset
{0<..<1}
proof -
  have 0 < left-seq n
  by (metis id-apply n-pos of-nat-eq-id of-nat-in-Nats order.asym
zero-lt-left-seq-lt-right-seq-both-pos)
  moreover have right-seq n < 1
  using right-seq-def2
  by (smt (verit, ccfv-SIG) inverse-1 inverse-le-imp-le mult-sign-intros(5)
n-pos of-nat-0-less-iff pi-gt3)
  ultimately show ?thesis
  using proper-interval by fastforce
qed
qed

```

```

show  $y \in \{\text{left-seq } n < \dots < \text{right-seq } n\}$ 
proof –
  have  $y \in \{\text{left-seq } n .. \text{right-seq } n\}$ 
    using  $y\text{-def}$  by blast
  moreover have  $y \neq \text{left-seq } n$ 
  proof(rule ccontr)
    assume  $\neg y \neq \text{left-seq } n$ 
    then have  $\text{deriv } f \ y \neq 0$ 
      using deriv-f first-and-second-order-conditions
      by (metis n-pos rel-simps(70) y-def)
    then show False
      by (simp add: y-def)
  qed
  moreover have  $y \neq \text{right-seq } n$ 
  proof(rule ccontr)
    assume  $\neg y \neq \text{right-seq } n$ 
    then have  $\text{deriv } f \ y \neq 0$ 
      using deriv-f first-and-second-order-conditions
      by (metis n-pos rel-simps(70) y-def)
    then show False
      by (simp add: y-def)
  qed
  ultimately show  $y \in \{\text{left-seq } n < \dots < \text{right-seq } n\}$ 
    by auto
  qed
show  $\text{deriv } f \ y = 0$  and  $0 < \text{deriv } (\text{deriv } f) \ y$ 
  using  $y\text{-def}$  by auto
show  $y \in \{\text{left-seq } n .. \text{right-seq } n\} \wedge \text{True}$ 
  using  $y\text{-def}$  by blast
qed
then show  $\exists y \in \{\text{left-seq } n .. \text{right-seq } n\}. \text{local-minimizer } f \ y$ 
  by blast
qed
show  $\exists x\text{-seq}. (\forall n. \text{local-minimizer } f \ (x\text{-seq } n) \wedge x\text{-seq } n \neq 0) \wedge x\text{-seq} \longrightarrow 0$ 
proof –
  define  $x\text{-seq}$  where
     $x\text{-seq } n = (\text{SOME } y. y \in \{\text{left-seq } (n+1) .. \text{right-seq } (n+1)\} \wedge \text{local-minimizer } f \ y)$ 
for  $n$ 
    have  $x\text{-seq-prop}: \forall n. x\text{-seq } n \in \{\text{left-seq } (n+1) .. \text{right-seq } (n+1)\} \wedge \text{local-minimizer } f \ (x\text{-seq } n)$ 
    by (metis (mono-tags, lifting) seq-of-local-minizers-exists someI-ex verit-eq-simplify(7) x-seq-def zero-eq-add-iff-both-eq-0)

    from  $x\text{-seq-prop}$  have  $\text{bounds}: \forall n. \text{left-seq } (n+1) \leq x\text{-seq } n \wedge x\text{-seq } n \leq \text{right-seq } (n+1)$ 
    by auto

  have  $\text{nonzero}: \forall n. x\text{-seq } n \neq 0$ 

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    by (metis Suc-eq-plus1 bounds id-apply nat.simps(3) not-less of-nat-eq-id
of-nat-in-Nats zero-lt-left-seq-lt-right-seq-both-pos)

  have left-seq-converges: left-seq  $\longrightarrow$  0
  proof (rule LIMSEQ-I)
    fix  $\varepsilon :: \text{real}$ 
    assume  $\varepsilon\text{-pos}$ :  $0 < \varepsilon$ 
    then obtain N where N-def:  $(N::\text{nat}) = \lceil 1 / (2 * \text{pi} * \varepsilon) \rceil + 1$ 
      by (metis add-mono-thms-linordered-field(5) arithmetic-simps(50) di-
vide-pos-pos
      mult-sign-intros(5) pi-gt-zero pos-int-cases semiring-norm(172)
      zero-less-ceiling zero-less-numeral)
    then have N-gt-0:  $N > 0$ 
    by (smt (verit)  $\varepsilon\text{-pos}$  divide-pos-pos gr0I int-ops(1) m2pi-less-pi mult-sign-intros(5)
zero-less-ceiling)

    have  $\forall n \geq N. |\text{left-seq } n| < \varepsilon$ 
    proof clarify
      fix  $n :: \text{nat}$ 
      assume n-ge:  $n \geq N$ 
      have left-seq-eqs:  $\text{left-seq } n = \text{inverse } ((5 * \text{pi} / 4) + 2 * \text{pi} * \text{real } n)$ 
      unfolding left-seq-def
      by (metis N-gt-0 id-apply left-seq-def linorder-not-less mult.commute n-ge
of-nat-eq-id of-nat-in-Nats vector-space-over-itself.scale-scale)
      show  $|\text{left-seq } n| < \varepsilon$ 
      proof -
        have  $|\text{left-seq } n| = 1 / ((5 * \text{pi} / 4) + 2 * \text{pi} * \text{real } n)$ 
        by (simp add: left-seq-eqs inverse-eq-divide)
        also have  $\dots \leq 1 / (2 * \text{pi} * \text{real } N)$ 
        by (smt (verit, best) N-gt-0 divide-nonneg-nonneg frac-le m2pi-less-pi
mult-left-mono mult-sign-intros(5) n-ge of-nat-0-less-iff of-nat-mono)
        also have  $\dots < 1 / (2 * \text{pi} * (\lceil 1 / (2 * \text{pi} * \varepsilon) \rceil))$ 
        by (smt (verit, best) N-def  $\varepsilon\text{-pos}$  ceiling-correct divide-pos-pos frac-less2
m2pi-less-pi mult-less-cancel-left-pos mult-sign-intros(5) of-int-1 of-int-add of-int-of-nat-eq)
        also have  $\dots \leq 1 / (2 * \text{pi} * (1 / (2 * \text{pi} * \varepsilon)))$ 
        by (smt (verit, ccfv-SIG)  $\varepsilon\text{-pos}$  ceiling-correct frac-le mult-left-mono
mult-sign-intros(5) pi-gt-zero zero-less-divide-iff)
        also have  $\dots = \varepsilon$ 
        by simp
      finally show ?thesis.
    qed
  qed
  then show  $\exists N. \forall n \geq N. \|\text{left-seq } n - 0\| < \varepsilon$ 
  by (metis cancel-comm-monoid-add-class.diff-zero real-norm-def)
qed
  have right-seq-converges: right-seq  $\longrightarrow$  0
  proof (rule LIMSEQ-I)
    fix  $\varepsilon :: \text{real}$ 
    assume  $\varepsilon\text{-pos}$ :  $0 < \varepsilon$ 

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then obtain  $N$  where  $N$ -def:  $(N::nat) = \lceil 1 / (2 * pi * \epsilon) \rceil + 1$ 
  by (metis add-mono-thms-linordered-field(5) arithmetic-simps(50) di-
vide-pos-pos
      mult-sign-intros(5) pi-gt-zero pos-int-cases semiring-norm(172)
      zero-less-ceiling zero-less-numeral)
hence  $N$ -gt-0:  $N > 0$ 
  by (smt (verit) eps-pos divide-pos-pos gr0I int-ops(1) m2pi-less-pi
mult-sign-intros(5)
      zero-less-ceiling)

have  $\forall n \geq N. |right\text{-seq } n| < \epsilon$ 
proof clarify
  fix  $n :: nat$ 
  assume  $n$ -ge:  $n \geq N$ 
  have right-seq-eqs:  $right\text{-seq } n = inverse (pi + 2 * pi * real n)$ 
    unfolding right-seq-def
  by (metis  $N$ -gt-0 id-apply linorder-not-less mult.commute mult.left-commute
 $n$ -ge of-nat-eq-id of-nat-in-Nats right-seq-def)
  show  $|right\text{-seq } n| < \epsilon$ 
  proof -
    have  $|right\text{-seq } n| = 1 / (pi + 2 * pi * real n)$ 
      by (simp add: right-seq-eqs inverse-eq-divide)
    also have  $\dots \leq 1 / (2 * pi * real N)$ 
      by (smt (verit, best)  $N$ -gt-0 divide-nonneg-nonneg frac-le m2pi-less-pi
mult-left-mono mult-sign-intros(5)  $n$ -ge of-nat-0-less-iff
of-nat-mono)
    also have  $\dots < 1 / (2 * pi * (\lceil 1 / (2 * pi * \epsilon) \rceil))$ 
      by (smt (verit, best)  $N$ -def eps-pos ceiling-correct divide-pos-pos frac-less2
m2pi-less-pi
mult-less-cancel-left-pos mult-sign-intros(5) of-int-1 of-int-add
of-int-of-nat-eq)
    also have  $\dots \leq 1 / (2 * pi * (1 / (2 * pi * \epsilon)))$ 
      by (smt (verit, ccfv-SIG) eps-pos ceiling-correct frac-le mult-left-mono
mult-sign-intros(5) pi-gt-zero zero-less-divide-iff)
    also have  $\dots = \epsilon$ 
      by simp
    finally show ?thesis
      by blast
  qed
qed
then show  $\exists no. \forall n \geq no. \|right\text{-seq } n - 0\| < \epsilon$ 
  by (metis cancel-comm-monoid-add-class.diff-zero real-norm-def)
qed
have  $x$ -seq-converges:  $x$ -seq  $\longrightarrow 0$ 
proof (rule LIMSEQ-I)
  fix  $\epsilon :: real$ 
  assume  $\epsilon$ -pos:  $0 < \epsilon$ 

obtain  $N_0$  where  $N_0$ :  $\forall n \geq N_0. \|left\text{-seq } (n+1) - 0\| < \epsilon$ 

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```

using left-seq-converges
by (meson LIMSEQ-iff  $\varepsilon$ -pos le-diff-conv)

obtain  $N_1$  where  $N_1: \forall n \geq N_1. \|right\text{-seq } (n+1) - 0\| < \varepsilon$ 
using right-seq-converges
by (meson LIMSEQ-iff  $\varepsilon$ -pos le-diff-conv)

obtain  $N$  where  $N = \max N_0 N_1$ 
by simp
hence  $N\text{-def}: N \geq N_0 \wedge N \geq N_1$ 
by simp

show  $\exists N. \forall n \geq N. \|x\text{-seq } n - 0\| < \varepsilon$ 
proof (intro exI[where  $x=N$ ] exI allI impI)
  fix  $n :: \text{nat}$ 
  assume  $N\text{-leq-}n: N \leq n$ 

  from bounds have  $left\text{-seq } (n+1) \leq x\text{-seq } n \wedge x\text{-seq } n \leq right\text{-seq } (n+1)$ 
  by auto
  hence  $\|x\text{-seq } n\| \leq \|left\text{-seq } (n+1)\| \vee \|x\text{-seq } n\| \leq \|right\text{-seq } (n+1)\|$ 
  by force
  moreover have  $\|left\text{-seq } (n+1)\| < \varepsilon \wedge \|right\text{-seq } (n+1)\| < \varepsilon$ 
  using  $N_0 N_1 N\text{-leq-}n N\text{-def}$  by auto
  ultimately show  $\|x\text{-seq } n - 0\| < \varepsilon$ 
  by auto
qed
qed
then show ?thesis
using nonzero x-seq-prop by blast
qed
qed
then show ?thesis
using zero-min f-cont not-isolated-minimizer-def strict-local-minimizer-at-0 by
auto
qed

end
theory Unconstrained-Optimization
imports Auxiliary-Facts
Minimizers-Definition
First-Order-Conditions
Second-Derivative-Test
Cont-Nonisolated-Strict-Local-Minimizer-Exists

begin

end

```

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