Transitive Union-Closed Families

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Abstract

We formalise a proof by Aaronson, Ellis and Leader showing that the Union-Closed Conjecture holds for the union-closed family generated by the cyclic translates of any fixed set.

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1 Transitive Union-Closed Families

A family of sets is union-closed if the union of any two sets from the family is in the family. The Union-Closed Conjecture is an open problem in combinatorics posed by Frankl in 1979. It states that for every finite, union-closed family of sets (other than the family containing only the empty set) there exists an element that belongs to at least half of the sets in the family. We formalise a proof by Aaronson, Ellis and Leader showing that the Union-Closed Conjecture holds for the union-closed family generated by the cyclic translates of any fixed set [1].

theory Transitive-Union-Closed-Families imports Pluennecke-Ruzsa-Inequality.Pluennecke-Ruzsa-Inequality

begin

no-notation equivalence.Partition (infix) // 75)

definition union-closed:: 'a set set \Rightarrow bool where union-closed $\mathcal{F} \equiv (\forall A \in \mathcal{F}, \forall B \in \mathcal{F}, A \cup B \in \mathcal{F})$ **abbreviation** set-difference :: $['a \ set, 'a \ set] \Rightarrow 'a \ set \ (infixl \setminus 65)$ where $A \setminus B \equiv A-B$

locale Family = additive-abelian-group + fixes Rassumes finG: finite Gassumes RG: $R \subseteq G$ assumes R-nonempty: $R \neq \{\}$

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definition union-closed-conjecture-property:: 'a set set \Rightarrow bool where union-closed-conjecture-property \mathcal{F} $\equiv \exists \mathcal{X} \subseteq \mathcal{F}. \exists x \in G. x \in \bigcap \mathcal{X} \land card \mathcal{X} \ge card \mathcal{F} / 2$

definition Neighbd $\equiv \lambda A$. sumset A R

definition Interior $\equiv \lambda A$. { $x \in G$. sumset {x} $R \subseteq A$ }

definition $\mathcal{F} \equiv Neighbd$ ' Pow G

We show that the family \mathcal{F} as defined above and appears in the statement of the theorem [1] is actually a finite, nonempty union-closed family indeed.

lemma card \mathcal{F} -gt0 [simp]: card $\mathcal{F} > 0$ and finite \mathcal{F} : finite \mathcal{F} using \mathcal{F} -def finG by fastforce+

 $\begin{array}{l} \textbf{lemma union-closed } \mathcal{F} \\ \textbf{proof-} \\ \textbf{have } *: \forall \ A \subseteq G. \ \forall \ B \subseteq G. \ (sumset \ A \ R) \cup (sumset \ B \ R) = sumset \ (A \cup B) \ R \\ \textbf{by } (simp \ add: \ sumset-subset-Un1) \\ \textbf{show } ?thesis \ \textbf{using } * \\ \textbf{by } (auto \ simp: \ union-closed-def \ \mathcal{F}-def \ Neighbd-def) \\ \textbf{qed} \end{array}$

lemma cardG-gt0: card G > 0using RG R-nonempty card-0-eq finG by blast

lemma \mathcal{F} -subset: $\mathcal{F} \subseteq Pow \ G$ **by** (simp add: Neighbd-def PowI \mathcal{F} -def image-subset-iff sumset-subset-carrier)

1.1 Proof of the main theorem

lemma card-Interior-le: **assumes** $S \subseteq G$ **shows** card (Interior S) \leq card S **proof** – **obtain** r where $r \in R$ **using** R-nonempty **by** blast **show** ?thesis

proof (*intro* card-inj-on-le) let $?f = (\lambda x. x \oplus r)$ **show** inj-on ?f (Interior S) ?f 'Interior $S \subseteq S$ using $RG \langle r \in R \rangle$ by (auto simp: Interior-def inj-on-def) show finite Susing assms finG finite-subset by blast qed qed **lemma** Interior-subset-G [iff]: Interior $S \subseteq G$ using Interior-def by auto **lemma** Neighbd-subset-G [iff]: Neighbd $S \subseteq G$ **by** (*simp add: Neighbd-def sumset-subset-carrier*) lemma average-ge: shows $(\sum S \in \mathcal{F}.(card S)) / card \mathcal{F} \ge card G / 2$ proofdefine f where $f \equiv \lambda S$. minusset $(G \setminus Interior S)$ The following corresponds to (1) in the paper. have 1: card S + card $(f S) \ge card G$ if $S \subseteq G$ for Sproofhave card (f S) = card G - card (Interior S)unfolding *f*-def by (metis Diff-subset Interior-subset-G card-Diff-subset card-minusset' finG *finite-subset*) with that show ?thesis using card-Interior-le by (metis (no-types, lifting) add.commute diff-le-mono2 le-diff-conv) qed The following corresponds to (2) in the paper. have 2: $f S = sumset (minusset (G \setminus S)) R$ if $S \subseteq G$ for S proofhave $*: x \in f S \longleftrightarrow x \in sumset (minusset (G \setminus S)) R$ if $x \in G$ for x proof – have $x \in f S \iff inverse \ x \notin Interior \ S$ using that minusset.simps by (fastforce simp: f-def)+ also have $\ldots \longleftrightarrow (sumset \{ inverse \ x \} \ R) \cap (G \setminus S) \neq \{ \}$ using sumset-subset-carrier that by (auto simp: Interior-def) also have $\ldots \longleftrightarrow x \in sumset \ (minusset \ (G \setminus S)) \ R$ proof assume L: sumset {inverse x} $R \cap (G \setminus S) \neq \{\}$ then obtain r where r: inverse $x \oplus r \notin S$ and $r \in R$ using $\langle S \subseteq G \rangle \langle x \in G \rangle$ by (auto simp: sumset-eq minusset-eq) then have inverse (inverse $x \oplus r$) \in minusset ($G \setminus S$) using RG that by auto moreover have x = inverse (inverse $x \oplus r$) $\oplus r$ using $RG \langle r \in R \rangle$ that commutative inverse-composition-commute invertible-right-inverse2

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by auto
       ultimately show x \in sumset (minusset (G \setminus S)) R
         by (metis RG \langle r \in R \rangle minusset-subset-carrier subset-eq sumset.simps)
     \mathbf{next}
       assume R: x \in sumset (minusset (G \setminus S)) R
       then obtain q r where *: q \in G q \notin S r \in R x = inverse q \oplus r
         by (metis Diff-iff minusset.simps sumset.cases)
       show sumset {inverse x} R \cap (G \setminus S) \neq \{\}
       proof
         assume sumset {inverse x} R \cap (G \setminus S) = \{\}
         then have g \notin sumset \{inverse \ x\} \ R
           using \langle g \notin S \rangle sumset-subset-carrier that by fastforce
         then have g \neq local.inverse \ (local.inverse \ g \oplus r) \oplus r
           using * RG that by (auto simp: sumset-eq)
         with * RG that show False
        by (metis commutative invertible invertible-left-inverse2 invertible-right-inverse2
subset-eq)
       qed
     qed
     finally show ?thesis .
   ged
   show ?thesis
   proof
     show f S \subseteq sumset (minusset (G \setminus S)) R
     using * f-def minusset-subset-carrier by blast
   \mathbf{next}
     show sumset (minusset (G \setminus S)) R \subseteq f S
     by (meson * subset-iff sumset-subset-carrier)
   qed
  qed
  then have f ` Pow \ G \subseteq \mathcal{F}
   by (auto simp: Neighbd-def \mathcal{F}-def minusset-subset-carrier)
    The following corresponds to (3) in the paper.
 have 3: Neighbd (Interior (sumset A R)) = sumset A R
   if A \subseteq G for A
   using that by (force simp: sumset-eq Neighbd-def Interior-def)
    "Putting everything together":
 moreover
 have sumset X R = sumset Y R
   \mathbf{if}\; X \subseteq \; G \; Y \subseteq \; G
      minusset (G \setminus Interior (sumset X R)) = minusset (G \setminus Interior (sumset Y R))
R))
   for X Y
   using that 3
  by (metis Diff-Diff-Int Int-absorb2 Interior-subset-G inf-commute minus-minusset)
  ultimately have inj-on f \mathcal{F}
```

 $\mathbf{by} \ (auto \ simp: \ inj\text{-}on\text{-}def \ \mathcal{F}\text{-}def \ f\text{-}def \ Neighbd\text{-}def)$

moreover have $f \, ` \, \mathcal{F} \subseteq \mathcal{F}$ using 2 \mathcal{F} -def $\langle f \ Pow \ G \subseteq \mathcal{F} \rangle$ by force moreover have $\mathcal{F} \subseteq f$ ' \mathcal{F} by (metis (inj-on $f \mathcal{F}$) ($f \mathcal{F} \subseteq \mathcal{F}$) endo-inj-surj finite \mathcal{F}) ultimately have bij-betw $f \mathcal{F} \mathcal{F}$ by (simp add: bij-betw-def) then have sum-card-eq: $(\sum S \in \mathcal{F}. card (f S)) = (\sum S \in \mathcal{F}. card S)$ **by** (*simp add: sum.reindex-bij-betw*) have card $G / 2 = (1 / (2 * card \mathcal{F})) * (\sum S \in \mathcal{F}. card G)$ by simp also have $\ldots \leq (1 / (2 * card \mathcal{F})) * (\sum S \in \mathcal{F}. card S + card (f S))$ by (intro sum-mono mult-left-mono of-nat-mono 1) (auto simp: \mathcal{F} -def) also have $\ldots = (1 / card \mathcal{F}) * (\sum S \in \mathcal{F}. card S)$ **by** (*simp add: sum-card-eq sum.distrib*) finally show ?thesis by argo qed

We have thus shown that the average size of a set in the family \mathcal{F} is at least |G|/2, proving the first part of Theorem 2 in the paper [1]. Using this, we will now show the main statement, i.e. that the Union-Closed Conjecture holds for the family \mathcal{F} .

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theorem Aaronson-Ellis-Leader-union-closed-conjecture:
shows union-closed-conjecture-property \mathcal{F}
proof -
```

— First, quite a big calculation not mentioned in the article: counting all the elements in two different ways.

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have *: (\sum S \in \mathcal{F}.(card S)) = (\sum x \in G. card \{S \in \mathcal{F}. x \in S\})
    using finite \mathcal{F} \mathcal{F}-subset
  proof induction
    case empty
    then show ?case
      by simp
  \mathbf{next}
    case (insert S \mathcal{G})
    then have A: \{T. (T = S \lor T \in \mathcal{G}) \land x \in T\}
                  = \{T \in \mathcal{G}. x \in T\} \cup (if x \in S then \{S\} else \{\})
      for x
      by auto
    have B: card {T. (T = S \lor T \in \mathcal{G}) \land x \in T}
            = card \{T \in \mathcal{G}. x \in T\} + (if x \in S then 1 else 0)
      for x
      by (simp add: A card-insert-if insert)
    have S = (\bigcup x \in G. if x \in S then \{x\} else \{\}
      using insert.prems by auto
    then have card S = card (\bigcup x \in G. if x \in S then \{x\} else \{\})
      by simp
    also have \ldots = (\sum i \in G. \ card \ (if \ i \in S \ then \ \{i\} \ else \ \{\}))
```

by (intro card-UN-disjoint) (auto simp: finG) also have $\ldots = (\sum x \in G. if x \in S then \ 1 else \ 0)$ **by** (force intro: sum.cong) finally have C: card $S = (\sum x \in G. \text{ if } x \in S \text{ then } 1 \text{ else } 0)$. show ?case using insert by (auto simp: sum.distrib B C) qed have $1/2 \leq (sum \ card \ \mathcal{F}) \ / \ (card \ \mathcal{F} * \ card \ G)$ using mult-right-mono [OF average-ge, of 1 / card G] using cardG-gt0 by (simp add: divide-simps split: if-splits) also have $\ldots = (\sum x \in G. ((card \{S \in \mathcal{F}. x \in S\}) / (card \mathcal{F}))) / card G$ **by** (*simp add*: * *sum-divide-distrib*) finally have **: $1/2 \leq (\sum x \in G. \ card \{S \in \mathcal{F}. \ x \in S\} \ / \ card \ \mathcal{F}) \ / \ card \ G$. - There is a typo in the paper (bottom of page): instead of $x \in S$ it says $x \in$ F. show ?thesis **proof** (*rule ccontr*) — Contradict the inequality proved above **assume** \neg union-closed-conjecture-property \mathcal{F} then have $A: \bigwedge \mathcal{X} x$. $[\mathcal{X} \subseteq \mathcal{F}; x \in G; x \in \bigcap \mathcal{X}] \Longrightarrow card \mathcal{X} < card \mathcal{F} / 2$ **by** (*fastforce simp: union-closed-conjecture-property-def*) have $(\sum x \in G. real (card \{S \in \mathcal{F}. x \in S\})) < (\sum x \in G. card \mathcal{F} / 2)$ **proof** (*intro sum-strict-mono*) fix x :: 'aassume $x \in G$ then have card $\{S \in \mathcal{F}. x \in S\} < card \mathcal{F} / 2$ by (intro A) auto then show real (card $\{S \in \mathcal{F}. x \in S\}$) < real (card \mathcal{F}) / 2 by blast qed (use unit-closed finG in auto) also have $\ldots = card \mathcal{F} * (card G / 2)$ by simp finally have B: $(\sum x \in G. real (card \{S \in \mathcal{F}. x \in S\})) < card \mathcal{F} * (card G / 2)$. have $(\sum x \in G. \text{ card } \{S \in \mathcal{F}. x \in S\} / \text{ card } \mathcal{F}) / \text{ card } G < 1/2$ using divide-strict-right-mono [OF B, of card $\mathcal{F} * card G$] using cardG- $qt\theta$ **by** (*simp add: divide-simps sum-divide-distrib*) with ****** show False by argo qed qed end

end

References

 J. Aaronson, D. Ellis, and I. Leader. A note on transitive union-closed families. 28(2), 2021. doihttps://doi.org/10.37236/9956.