

The Theorem of Three Circles

Fox Thomson, Wenda Li

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Abstract

The Descartes test based on Bernstein coefficients and Descartes' rule of signs effectively (over-)approximates the number of real roots of a univariate polynomial over an interval. In this entry we formalise the theorem of three circles (Theorem 10.50 in [1]), which gives sufficient conditions for when the Descartes test returns 0 or 1. This is the first step for efficient root isolation.

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1 Misc results about polynomials

theory *RRI-Misc* **imports**

HOL-Computational-Algebra.Computational-Algebra

Budan-Fourier.BF-Misc

Polynomial-Interpolation.Ring-Hom-Poly

begin

1.1 Misc

declare *pcompose-pCons*[*simp del*]

lemma *Setcompr-subset*: $\bigwedge f P S. \{f x \mid x. P x\} \subseteq S = (\forall x. P x \longrightarrow f x \in S)$
<proof>

lemma *map-cong'*:

assumes *xs = map h ys* **and** $\bigwedge y. y \in \text{set } ys \implies f (h y) = g y$

shows *map f xs = map g ys*

<proof>

lemma *nth-default-replicate-eq*:

nth-default dflt (replicate n x) i = (if i < n then x else dflt)

<proof>

lemma *square-bounded-less*:

fixes *a b::'a :: linordered-ring-strict*

shows $-a < b \wedge b < a \implies b*b < a*a$

<proof>

lemma *square-bounded-le*:

fixes *a b::'a :: linordered-ring-strict*

shows $-a \leq b \wedge b \leq a \implies b*b \leq a*a$

<proof>

context *vector-space*

begin

lemma *card-le-dim-spanning*:

assumes *BV: B ⊆ V*

and *VB: V ⊆ span B*

and *fB: finite B*

and *dVB: dim V ≥ card B*

shows *independent B*

<proof>

end

1.2 Misc results about polynomials

lemma *smult-power*: $smult (x \hat{n}) (p \hat{n}) = smult x p \hat{n}$
 ⟨proof⟩

lemma *reflect-poly-monom*: $reflect-poly (monom n i) = monom n 0$
 ⟨proof⟩

lemma *poly-eq-by-eval*:
 fixes $P Q :: 'a :: \{comm-ring-1, ring-no-zero-divisors, ring-char-0\}$ poly
 assumes $h: \bigwedge x. poly P x = poly Q x$ shows $P = Q$
 ⟨proof⟩

lemma *poly-binomial*:
 $[(1::'a::comm-ring-1), 1] \hat{n} = (\sum k \leq n. monom (of-nat (n choose k)) k)$
 ⟨proof⟩

lemma *degree-0-iff*: $degree P = 0 \longleftrightarrow (\exists a. P = [:a:])$
 ⟨proof⟩

interpretation *poly-vs*: *vector-space smult*
 ⟨proof⟩

lemma *degree-subspace*: $poly-vs.subspace \{x. degree x \leq n\}$
 ⟨proof⟩

lemma *monom-span*:
 $poly-vs.span \{monom 1 x \mid x. x \leq p\} = \{(x::'a::field poly). degree x \leq p\}$
 (is ?L = ?R)
 ⟨proof⟩

lemma *monom-independent*:
 $poly-vs.independent \{monom (1::'a::field) x \mid x. x \leq p\}$
 ⟨proof⟩

lemma *dim-degree*: $poly-vs.dim \{x. degree x \leq n\} = n + 1$
 ⟨proof⟩

lemma *degree-div*:
 fixes $p q :: ('a::idom-divide) poly$
 assumes $q \text{ dvd } p$
 shows $degree (p \text{ div } q) = degree p - degree q$ ⟨proof⟩

lemma *lead-coeff-div*:
 fixes $p q :: ('a::\{idom-divide, inverse\}) poly$
 assumes $q \text{ dvd } p$
 shows $lead-coeff (p \text{ div } q) = lead-coeff p / lead-coeff q$ ⟨proof⟩

lemma *complex-poly-eq*:
 $r = map-poly \text{ of-real } (map-poly Re r) + smult i (map-poly \text{ of-real } (map-poly Im$

r))
 ⟨proof⟩

lemma *complex-poly-cong*:
 (map-poly Re p = map-poly Re q ∧ map-poly Im p = map-poly Im q) = (p = q)
 ⟨proof⟩

lemma *map-poly-Im-of-real*: map-poly Im (map-poly of-real p) = 0
 ⟨proof⟩

lemma *mult-map-poly-imp-map-poly*:
 assumes map-poly complex-of-real q = r * map-poly complex-of-real p
 p ≠ 0
 shows r = map-poly complex-of-real (map-poly Re r)
 ⟨proof⟩

lemma *map-poly-dvd*:
 fixes p q::real poly
 assumes hdvd: map-poly complex-of-real p dvd
 map-poly complex-of-real q q ≠ 0
 shows p dvd q
 ⟨proof⟩

lemma *div-poly-eq-0*:
 fixes p q:(*'a::idom-divide*) poly
 assumes q dvd p poly (p div q) x = 0 q ≠ 0
 shows poly p x = 0
 ⟨proof⟩

lemma *poly-map-poly-of-real-cnj*:
 poly (map-poly of-real p) (cnj z) = cnj (poly (map-poly of-real p) z)
 ⟨proof⟩

An induction rule on real polynomials, if $P \neq 0$ then either $(X - x)|P$ or $(X - z)(X - cnjz)|P$, we induct by dividing by these polynomials.

lemma *real-poly-roots-induct*:
 fixes P::real poly ⇒ bool and p::real poly
 assumes IH-real: $\bigwedge p x. P p \implies P (p * [:-x, 1:])$
 and IH-complex: $\bigwedge p a b. b \neq 0 \implies P p$
 $\implies P (p * [: a*a + b*b, -2*a, 1 :])$
 and H0: $\bigwedge a. P [:a:]$
 defines d ≡ degree p
 shows P p
 ⟨proof⟩

1.3 The reciprocal polynomial

definition *reciprocal-poly* :: nat ⇒ *'a::zero poly* ⇒ *'a poly*
 where *reciprocal-poly* p P =

$Poly (rev ((coeffs P) @ (replicate (p - degree P) 0)))$

lemma *reciprocal-0*: *reciprocal-poly p 0 = 0* $\langle proof \rangle$

lemma *reciprocal-1*: *reciprocal-poly p 1 = monom 1 p*
 $\langle proof \rangle$

lemma *coeff-reciprocal*:
assumes *hi*: $i \leq p$ and *hP*: $degree P \leq p$
shows $coeff (reciprocal-poly p P) i = coeff P (p - i)$
 $\langle proof \rangle$

lemma *coeff-reciprocal-less*:
assumes *hn*: $p < i$ and *hP*: $degree P \leq p$
shows $coeff (reciprocal-poly p P) i = 0$
 $\langle proof \rangle$

lemma *reciprocal-monom*:
assumes $n \leq p$
shows $reciprocal-poly p (monom a n) = monom a (p-n)$
 $\langle proof \rangle$

lemma *reciprocal-degree*: *reciprocal-poly (degree P) P = reflect-poly P*
 $\langle proof \rangle$

lemma *degree-reciprocal*:
fixes $P :: ('a::zero) poly$
assumes *hP*: $degree P \leq p$
shows $degree (reciprocal-poly p P) \leq p$
 $\langle proof \rangle$

lemma *reciprocal-0-iff*:
assumes *hP*: $degree P \leq p$
shows $(reciprocal-poly p P = 0) = (P = 0)$
 $\langle proof \rangle$

lemma *poly-reciprocal*:
fixes $P :: 'a::field poly$
assumes *hp*: $degree P \leq p$ and *hx*: $x \neq 0$
shows $poly (reciprocal-poly p P) x = x^p * (poly P (inverse x))$
 $\langle proof \rangle$

lemma *reciprocal-fcompose*:
fixes $P :: ('a::\{ring-char-0,field\}) poly$
assumes *hP*: $degree P \leq p$
shows $reciprocal-poly p P = monom 1 (p - degree P) * fcompose P 1 [:0, 1:]$
 $\langle proof \rangle$

lemma *reciprocal-reciprocal*:

fixes $P :: 'a::\{\text{field}, \text{ring-char-0}\}$ poly
assumes $hP: \text{degree } P \leq p$
shows $\text{reciprocal-poly } p (\text{reciprocal-poly } p P) = P$
 <proof>

lemma *reciprocal-smult*:
fixes $P :: 'a::\text{idom}$ poly
assumes $h: \text{degree } P \leq p$
shows $\text{reciprocal-poly } p (\text{smult } n P) = \text{smult } n (\text{reciprocal-poly } p P)$
 <proof>

lemma *reciprocal-add*:
fixes $P Q :: 'a::\text{comm-semiring-0}$ poly
assumes $\text{degree } P \leq p$ **and** $\text{degree } Q \leq p$
shows $\text{reciprocal-poly } p (P + Q) = \text{reciprocal-poly } p P + \text{reciprocal-poly } p Q$
 (is ?L = ?R)
 <proof>

lemma *reciprocal-diff*:
fixes $P Q :: 'a::\text{comm-ring}$ poly
assumes $\text{degree } P \leq p$ **and** $\text{degree } Q \leq p$
shows $\text{reciprocal-poly } p (P - Q) = \text{reciprocal-poly } p P - \text{reciprocal-poly } p Q$
 <proof>

lemma *reciprocal-sum*:
fixes $P :: 'a \Rightarrow 'b::\text{comm-semiring-0}$ poly
assumes $hP: \bigwedge k. \text{degree } (P k) \leq p$
shows $\text{reciprocal-poly } p (\sum k \in A. P k) = (\sum k \in A. \text{reciprocal-poly } p (P k))$
 <proof>

lemma *reciprocal-mult*:
fixes $P Q :: 'a::\{\text{ring-char-0}, \text{field}\}$ poly
assumes $\text{degree } (P * Q) \leq p$
and $\text{degree } P \leq p$ **and** $\text{degree } Q \leq p$
shows $\text{monom } 1 p * \text{reciprocal-poly } p (P * Q) =$
 $\text{reciprocal-poly } p P * \text{reciprocal-poly } p Q$
 <proof>

lemma *reciprocal-reflect-poly*:
fixes $P :: 'a::\{\text{ring-char-0}, \text{field}\}$ poly
assumes $hP: \text{degree } P \leq p$
shows $\text{reciprocal-poly } p P = \text{monom } 1 (p - \text{degree } P) * \text{reflect-poly } P$
 <proof>

lemma *map-poly-reciprocal*:
assumes $\text{degree } P \leq p$ **and** $f 0 = 0$
shows $\text{map-poly } f (\text{reciprocal-poly } p P) = \text{reciprocal-poly } p (\text{map-poly } f P)$
 <proof>

1.4 More about *proots-count*

lemma *proots-count-monom*:

assumes $0 \notin A$

shows $\text{proots-count } (\text{monom } 1 \ d) \ A = 0$

<proof>

lemma *proots-count-reciprocal*:

fixes $P::'a::\{\text{ring-char-0,field}\}$ *poly*

assumes hP : $\text{degree } P \leq p$ **and** $h0$: $P \neq 0$ **and** $h0'$: $0 \notin A$

shows $\text{proots-count } (\text{reciprocal-poly } p \ P) \ A = \text{proots-count } P \ \{x. \text{inverse } x \in A\}$

<proof>

lemma *proots-count-reciprocal'*:

fixes $P::\text{real } \text{poly}$

assumes hP : $\text{degree } P \leq p$ **and** $h0$: $P \neq 0$

shows $\text{proots-count } P \ \{x. 0 < x \wedge x < 1\} =$

$\text{proots-count } (\text{reciprocal-poly } p \ P) \ \{x. 1 < x\}$

<proof>

lemma *proots-count-pos*:

assumes $\text{proots-count } P \ S > 0$

shows $\exists x \in S. \text{poly } P \ x = 0$

<proof>

lemma *proots-count-of-root-set*:

assumes $P \neq 0$ $R \subseteq S$ **and** $\bigwedge x. x \in R \implies \text{poly } P \ x = 0$

shows $\text{proots-count } P \ S \geq \text{card } R$

<proof>

lemma *proots-count-of-root*: **assumes** $P \neq 0$ $x \in S$ $\text{poly } P \ x = 0$

shows $\text{proots-count } P \ S > 0$

<proof>

1.5 More about *changes*

lemma *changes-nonneg*: $0 \leq \text{changes } xs$

<proof>

lemma *changes-replicate-0*: **shows** $\text{changes } (\text{replicate } n \ 0) = 0$

<proof>

lemma *changes-append-replicate-0*: $\text{changes } (xs \ @ \ \text{replicate } n \ 0) = \text{changes } xs$

<proof>

lemma *changes-scale-Cons*:

fixes $xs::\text{real list}$ **assumes** hs : $s > 0$

shows $\text{changes } (s * x \ \# \ xs) = \text{changes } (x \ \# \ xs)$

<proof>

lemma *changes-scale*:

fixes $xs::('a::\text{linordered-idom}) \text{ list}$

assumes $hs: \bigwedge i. i < n \implies s \ i > 0$ **and** $hn: \text{length } xs \leq n$

shows $\text{changes } [s \ i * (\text{nth-default } 0 \ xs \ i). \ i \leftarrow [0..<n]] = \text{changes } xs$

<proof>

lemma *changes-scale-const*: **fixes** $xs::'a::\text{linordered-idom} \text{ list}$

assumes $hs: s \neq 0$

shows $\text{changes } (\text{map } ((* \ s) \ xs) = \text{changes } xs$

<proof>

lemma *changes-snoc*: **fixes** $xs::'a::\text{linordered-idom} \text{ list}$

shows $\text{changes } (xs \ @ \ [b, a]) = (\text{if } a * b < 0 \ \text{then } 1 + \text{changes } (xs \ @ \ [b])$
 $\text{else if } b = 0 \ \text{then } \text{changes } (xs \ @ \ [a]) \ \text{else } \text{changes } (xs \ @ \ [b]))$

<proof>

lemma *changes-rev*: **fixes** $xs::'a::\text{linordered-idom} \text{ list}$

shows $\text{changes } (\text{rev } xs) = \text{changes } xs$

<proof>

lemma *changes-rev-about*: **fixes** $xs::'a::\text{linordered-idom} \text{ list}$

shows $\text{changes } (\text{replicate } (p - \text{length } xs) \ 0 \ @ \ \text{rev } xs) = \text{changes } xs$

<proof>

lemma *changes-add-between*:

assumes $a \leq x$ **and** $x \leq b$

shows $\text{changes } (as \ @ \ [a, b] \ @ \ bs) = \text{changes } (as \ @ \ [a, x, b] \ @ \ bs)$

<proof>

lemma *changes-all-nonneg*: **assumes** $\bigwedge i. \text{nth-default } 0 \ xs \ i \geq 0$ **shows** $\text{changes } xs = 0$

<proof>

lemma *changes-pCons*: $\text{changes } (\text{coeffs } (pCons \ 0 \ f)) = \text{changes } (\text{coeffs } f)$

<proof>

lemma *changes-increasing*:

assumes $\bigwedge i. i < \text{length } xs - 1 \implies xs \ ! \ (i + 1) \geq xs \ ! \ i$

and $\text{length } xs > 1$

and $\text{hd } xs < 0$

and $\text{last } xs > 0$

shows $\text{changes } xs = 1$

<proof>

end

2 Bernstein Polynomials over the interval $[0, 1]$

theory *Bernstein-01*

imports *HOL-Computational-Algebra.Computational-Algebra*
Budan-Fourier.Budan-Fourier
RRI-Misc

begin

The theorem of three circles is a statement about the Bernstein coefficients of a polynomial, the coefficients when a polynomial is expressed as a sum of Bernstein polynomials. These coefficients behave nicely under translations and rescaling and are the coefficients of a particular polynomial in the $[0, 1]$ case. We shall define the $[0, 1]$ case now and consider the general case later, deriving all the results by rescaling.

2.1 Definition and basic results

definition *Bernstein-Poly-01* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real poly}$ **where**
Bernstein-Poly-01 j $p = (\text{monom } (p \text{ choose } j) \ j)$
 $\quad * (\text{monom } 1 \ (p-j) \ \circ_p \ [:1, -1:])$

lemma *degree-Bernstein*:

assumes $hb: j \leq p$

shows $\text{degree } (\text{Bernstein-Poly-01 } j \ p) = p$

$\langle \text{proof} \rangle$

lemma *coeff-gt*:

assumes $hb: j > p$

shows $\text{Bernstein-Poly-01 } j \ p = 0$

$\langle \text{proof} \rangle$

lemma *degree-Bernstein-le*: $\text{degree } (\text{Bernstein-Poly-01 } j \ p) \leq p$

$\langle \text{proof} \rangle$

lemma *poly-Bernstein-nonneg*:

assumes $x \geq 0$ **and** $1 \geq x$

shows $\text{poly } (\text{Bernstein-Poly-01 } j \ p) \ x \geq 0$

$\langle \text{proof} \rangle$

lemma *Bernstein-symmetry*:

assumes $j \leq p$

shows $(\text{Bernstein-Poly-01 } j \ p) \ \circ_p \ [:1, -1:] = \text{Bernstein-Poly-01 } (p-j) \ p$

$\langle \text{proof} \rangle$

2.2 Bernstein-Poly-01 and reciprocal-poly

lemma *Bernstein-reciprocal*:

$\text{reciprocal-poly } p \ (\text{Bernstein-Poly-01 } i \ p)$

$= \text{smult } (p \text{ choose } i) \ ([: -1, 1:] \ \frown (p-i))$

$\langle \text{proof} \rangle$

lemma *Bernstein-reciprocal-translate*:

reciprocal-poly p (*Bernstein-Poly-01* i p) \circ_p $[:1, 1:] =$
monom (p choose i) $(p - i)$
 \langle *proof* \rangle

lemma *coeff-Bernstein-sum-01*: **fixes** $b::\text{nat} \Rightarrow \text{real}$ **assumes** $hi: p \geq i$
shows
coeff (*reciprocal-poly* p
 $(\sum x = 0..p. \text{smult } (b \ x) \ (\text{Bernstein-Poly-01 } x \ p)) \circ_p$ $[:1, 1:]$)
 $(p - i) = (p \ \text{choose } i) * (b \ i)$ (**is** $?L = ?R$)
 \langle *proof* \rangle

lemma *Bernstein-sum-01*: **assumes** $hP: \text{degree } P \leq p$
shows
 $P = (\sum j = 0..p. \text{smult}$
 $(\text{inverse } (\text{real } (p \ \text{choose } j))) * \text{coeff } (\text{reciprocal-poly } p \ P \circ_p$ $[:1, 1:]$) $(p-j)$)
 $(\text{Bernstein-Poly-01 } j \ p))$
 \langle *proof* \rangle

lemma *Bernstein-Poly-01-span1*:
assumes $hP: \text{degree } P \leq p$
shows $P \in \text{poly-vs.span } \{\text{Bernstein-Poly-01 } x \ p \mid x. x \leq p\}$
 \langle *proof* \rangle

lemma *Bernstein-Poly-01-span*:
 $\text{poly-vs.span } \{\text{Bernstein-Poly-01 } x \ p \mid x. x \leq p\}$
 $= \{x. \text{degree } x \leq p\}$
 \langle *proof* \rangle

2.3 Bernstein coefficients and changes

definition *Bernstein-coeffs-01* $:: \text{nat} \Rightarrow \text{real poly} \Rightarrow \text{real list}$ **where**
Bernstein-coeffs-01 $p \ P =$
 $[(\text{inverse } (\text{real } (p \ \text{choose } j))) * \text{coeff } (\text{reciprocal-poly } p \ P \circ_p$ $[:1, 1:]$) $(p-j)]. j \leftarrow [0..<(p+1)]$

lemma *length-Bernstein-coeffs-01*: $\text{length } (\text{Bernstein-coeffs-01 } p \ P) = p + 1$
 \langle *proof* \rangle

lemma *nth-default-Bernstein-coeffs-01*: **assumes** $\text{degree } P \leq p$
shows $\text{nth-default } 0 \ (\text{Bernstein-coeffs-01 } p \ P) \ i =$
 $\text{inverse } (p \ \text{choose } i) * \text{coeff } (\text{reciprocal-poly } p \ P \circ_p$ $[:1, 1:]$) $(p-i)$
 \langle *proof* \rangle

lemma *Bernstein-coeffs-01-sum*: **assumes** $\text{degree } P \leq p$
shows $P = (\sum j = 0..p. \text{smult } (\text{nth-default } 0 \ (\text{Bernstein-coeffs-01 } p \ P) \ j)$
 $(\text{Bernstein-Poly-01 } j \ p))$
 \langle *proof* \rangle

definition *Bernstein-changes-01* :: *nat* \Rightarrow *real poly* \Rightarrow *int* **where**
Bernstein-changes-01 *p P* = *nat* (*changes* (*Bernstein-coeffs-01* *p P*))

lemma *Bernstein-changes-01-def'*:
Bernstein-changes-01 *p P* = *nat* (*changes* [(*inverse* (*real* (*p choose j*)) *
coeff (*reciprocal-poly* *p P* \circ_p [:1, 1:] (*p-j*)). *j* \leftarrow [0..*p* + 1]])
 \langle *proof* \rangle

lemma *Bernstein-changes-01-eq-changes*:
assumes *hP*: *degree P* \leq *p*
shows *Bernstein-changes-01* *p P* =
changes (*coeffs* ((*reciprocal-poly* *p P*) \circ_p [:1, 1:])))
 \langle *proof* \rangle

lemma *Bernstein-changes-01-test*: **fixes** *P*::*real poly*
assumes *hP*: *degree P* \leq *p* **and** *h0*: *P* \neq 0
shows *roots-count* *P* {*x*. 0 < *x* \wedge *x* < 1} \leq *Bernstein-changes-01* *p P* \wedge
even (*Bernstein-changes-01* *p P* - *roots-count* *P* {*x*. 0 < *x* \wedge *x* < 1})
 \langle *proof* \rangle

2.4 Expression as a Bernstein sum

lemma *Bernstein-coeffs-01-0*: *Bernstein-coeffs-01* *p* 0 = *replicate* (*p*+1) 0
 \langle *proof* \rangle

lemma *Bernstein-coeffs-01-1*: *Bernstein-coeffs-01* *p* 1 = *replicate* (*p*+1) 1
 \langle *proof* \rangle

lemma *Bernstein-coeffs-01-x*: **assumes** *p* \neq 0
shows *Bernstein-coeffs-01* *p* (*monom* 1 1) = [*i*/*p*. *i* \leftarrow [0..*p*+1]]
 \langle *proof* \rangle

lemma *Bernstein-coeffs-01-add*:
assumes *degree P* \leq *p* **and** *degree Q* \leq *p*
shows *nth-default* 0 (*Bernstein-coeffs-01* *p* (*P* + *Q*)) *i* =
nth-default 0 (*Bernstein-coeffs-01* *p* *P*) *i* +
nth-default 0 (*Bernstein-coeffs-01* *p* *Q*) *i*
 \langle *proof* \rangle

lemma *Bernstein-coeffs-01-smult*:
assumes *degree P* \leq *p*
shows *nth-default* 0 (*Bernstein-coeffs-01* *p* (*smult* *a* *P*)) *i* =
a * *nth-default* 0 (*Bernstein-coeffs-01* *p* *P*) *i*
 \langle *proof* \rangle

end

3 Bernstein Polynomials over any finite interval

theory *Bernstein*
imports *Bernstein-01*
begin

3.1 Definition and relation to Bernstein Polynomials over $[0, 1]$

definition *Bernstein-Poly* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real poly}$ **where**
Bernstein-Poly j p c $d = \text{smult } ((p \text{ choose } j)/(d - c) \widehat{p})$
 $((\text{monom } 1 \ j) \circ_p [-c, 1:]) * (\text{monom } 1 \ (p-j) \circ_p [:d, -1:])))$

lemma *Bernstein-Poly-altdef*:
assumes $c \neq d$ **and** $j \leq p$
shows *Bernstein-Poly* j p c $d = \text{smult } (p \text{ choose } j)$
 $([:-c/(d-c), 1/(d-c):] \widehat{j} * [:d/(d-c), -1/(d-c):] \widehat{(p-j)})$
(is ?L = ?R)
 $\langle \text{proof} \rangle$

lemma *Bernstein-Poly-nonneg*:
assumes $c \leq x$ **and** $x \leq d$
shows *poly* (*Bernstein-Poly* j p c d) $x \geq 0$
 $\langle \text{proof} \rangle$

lemma *Bernstein-Poly-01*: *Bernstein-Poly* j p 0 $1 = \text{Bernstein-Poly-01 } j$ p
 $\langle \text{proof} \rangle$

lemma *Bernstein-Poly-rescale*:
assumes $a \neq b$
shows *Bernstein-Poly* j p c $d \circ_p [:a, 1:] \circ_p [:0, b-a:]$
 $= \text{Bernstein-Poly } j$ p $((c-a)/(b-a))$ $((d-a)/(b-a))$
(is ?L = ?R)
 $\langle \text{proof} \rangle$

lemma *Bernstein-Poly-rescale-01*:
assumes $c \neq d$
shows *Bernstein-Poly* j p c $d \circ_p [:c, 1:] \circ_p [:0, d-c:]$
 $= \text{Bernstein-Poly-01 } j$ p
 $\langle \text{proof} \rangle$

lemma *Bernstein-Poly-eq-rescale-01*:
assumes $c \neq d$
shows *Bernstein-Poly* j p c $d = \text{Bernstein-Poly-01 } j$ p
 $\circ_p [:0, 1/(d-c):] \circ_p [-c, 1:]$
 $\langle \text{proof} \rangle$

lemma *coeff-Bernstein-sum*:
fixes $b::\text{nat} \Rightarrow \text{real}$ **and** $p::\text{nat}$ **and** c $d::\text{real}$

defines $P \equiv (\sum j = 0..p. (smult (b j) (Bernstein-Poly j p c d)))$
assumes $i \leq p$ **and** $c \neq d$
shows $coeff ((reciprocal-poly p (P \circ_p [:c, 1:]$
 $\circ_p [:0, d-c:])) \circ_p [:1, 1:]) (p - i) = (p \text{ choose } i) * (b i)$
 $\langle proof \rangle$

lemma *Bernstein-sum*:

assumes $c \neq d$ **and** $degree P \leq p$
shows $P = (\sum j = 0..p. smult (inverse (real (p \text{ choose } j))$
 $* coeff (reciprocal-poly p (P \circ_p [:c, 1:] \circ_p [:0, d-c:]$
 $\circ_p [:1, 1:]) (p-j)) (Bernstein-Poly j p c d))$
 $\langle proof \rangle$

lemma *Bernstein-Poly-span1*:

assumes $c \neq d$ **and** $degree P \leq p$
shows $P \in poly\text{-vs.}\text{span} \{Bernstein-Poly x p c d \mid x. x \leq p\}$
 $\langle proof \rangle$

lemma *Bernstein-Poly-span*:

assumes $c \neq d$
shows $poly\text{-vs.}\text{span} \{Bernstein-Poly x p c d \mid x. x \leq p\} = \{x. degree x \leq p\}$
 $\langle proof \rangle$

lemma *Bernstein-Poly-independent*: **assumes** $c \neq d$

shows $poly\text{-vs.}\text{independent} \{Bernstein-Poly x p c d \mid x. x \in \{..p\}\}$
 $\langle proof \rangle$

3.2 Bernstein coefficients and changes over any interval

definition *Bernstein-coeffs* ::

$nat \Rightarrow real \Rightarrow real \Rightarrow real \text{ poly} \Rightarrow real \text{ list}$ **where**
Bernstein-coeffs $p c d P =$
 $[(inverse (real (p \text{ choose } j)) *$
 $coeff (reciprocal-poly p (P \circ_p [:c, 1:] \circ_p [:0, d-c:] \circ_p [:1, 1:]) (p-j)).$
 $j \leftarrow [0..(p+1)]]$

lemma *Bernstein-coeffs-eq-rescale*: **assumes** $c \neq d$

shows $Bernstein-coeffs p c d P = Bernstein-coeffs-01 p (P \circ_p [:c, 1:] \circ_p [:0,$
 $d-c:])$
 $\langle proof \rangle$

lemma *nth-default-Bernstein-coeffs*: **assumes** $degree P \leq p$

shows $nth\text{-default } 0 (Bernstein-coeffs p c d P) i =$
 $inverse (p \text{ choose } i) * coeff$
 $(reciprocal-poly p (P \circ_p [:c, 1:] \circ_p [:0, d-c:] \circ_p [:1, 1:]) (p-i))$
 $\langle proof \rangle$

lemma *Bernstein-coeffs-sum*: **assumes** $c \neq d$ **and** hP : $degree P \leq p$

shows $P = (\sum j = 0..p. smult (nth\text{-default } 0 (Bernstein-coeffs p c d P) j)$

(Bernstein-Poly j p c d)
 ⟨proof⟩

definition *Bernstein-changes* :: nat ⇒ real ⇒ real ⇒ real poly ⇒ int **where**
Bernstein-changes p c d P = nat (changes (Bernstein-coeffs p c d P))

lemma *Bernstein-changes-eq-rescale*: **assumes** $c \neq d$ **and** degree $P \leq p$
shows *Bernstein-changes* p c d P =
Bernstein-changes-01 p (P ◦_p [:c, 1:] ◦_p [:0, d-c:])
 ⟨proof⟩

This is related and mostly equivalent to previous Descartes test [3]

lemma *Bernstein-changes-test*:
fixes P::real poly
assumes degree $P \leq p$ **and** $P \neq 0$ **and** $c < d$
shows *proots-count* P {x. $c < x \wedge x < d$ } ≤ *Bernstein-changes* p c d P ∧
 even (Bernstein-changes p c d P − *proots-count* P {x. $c < x \wedge x < d$ })
 ⟨proof⟩

3.3 The control polygon of a polynomial

definition *control-points* ::
 nat ⇒ real ⇒ real ⇒ real poly ⇒ (real × real) list
where
control-points p c d P =
 [(((real i)*d + (real (p - i))*c)/p,
 nth-default 0 (Bernstein-coeffs p c d P) i).
 i ← [0.. $(p+1)$]]

lemma *line-above*:
fixes a b c d :: real **and** p :: nat **and** P :: real poly
assumes *hline*: $\bigwedge i. i \leq p \implies a * (((real i)*d + (real (p - i))*c)/p) + b \geq$
 nth-default 0 (Bernstein-coeffs p c d P) i
and *hp*: $p \neq 0$ **and** *hcd*: $c \neq d$ **and** *hP*: degree $P \leq p$
shows $\bigwedge x. c \leq x \implies x \leq d \implies a*x + b \geq \text{poly } P x$
 ⟨proof⟩

end

4 Normal Polynomials

theory *Normal-Poly*
imports *RRI-Misc*
begin

Here we define normal polynomials as defined in Basu, S., Pollack, R., Roy, M.-F.: Algorithms in Real Algebraic Geometry. Springer Berlin Heidelberg, Berlin, Heidelberg (2016).

definition *normal-poly* :: ('a::{comm-ring-1,ord}) poly \Rightarrow bool **where**
normal-poly p \equiv

(p \neq 0) \wedge
 $(\forall i. 0 \leq \text{coeff } p \ i) \wedge$
 $(\forall i. \text{coeff } p \ i * \text{coeff } p \ (i+2) \leq (\text{coeff } p \ (i+1))^2) \wedge$
 $(\forall i \ j \ k. i \leq j \longrightarrow j \leq k \longrightarrow 0 < \text{coeff } p \ i$
 $\longrightarrow 0 < \text{coeff } p \ k \longrightarrow 0 < \text{coeff } p \ j)$

lemma *normal-non-zero*: *normal-poly* p \Longrightarrow p \neq 0

<proof>

lemma *normal-coeff-nonneg*: *normal-poly* p \Longrightarrow 0 \leq coeff p i

<proof>

lemma *normal-poly-coeff-mult*:

normal-poly p \Longrightarrow coeff p i * coeff p (i+2) \leq (coeff p (i+1))^2

<proof>

lemma *normal-poly-pos-interval*:

normal-poly p \Longrightarrow i \leq j \Longrightarrow j \leq k \Longrightarrow 0 < coeff p i \Longrightarrow 0 < coeff p k
 \Longrightarrow 0 < coeff p j

<proof>

lemma *normal-polyI*:

assumes (p \neq 0)

and ($\bigwedge i. 0 \leq \text{coeff } p \ i$)

and ($\bigwedge i. \text{coeff } p \ i * \text{coeff } p \ (i+2) \leq (\text{coeff } p \ (i+1))^2$)

and ($\bigwedge i \ j \ k. i \leq j \Longrightarrow j \leq k \Longrightarrow 0 < \text{coeff } p \ i \Longrightarrow 0 < \text{coeff } p \ k \Longrightarrow 0 <$

coeff p j)

shows *normal-poly* p

<proof>

lemma *linear-normal-iff*:

fixes x::real

shows *normal-poly* [:-x, 1:] \longleftrightarrow x \leq 0

<proof>

lemma *quadratic-normal-iff*:

fixes z::complex

shows *normal-poly* [(cmod z)^2, -2*Re z, 1:]

\longleftrightarrow Re z \leq 0 \wedge 4*(Re z)^2 \geq (cmod z)^2

<proof>

lemma *normal-of-no-zero-root*:

fixes f::real poly

assumes hzero: poly f 0 \neq 0 **and** hdeg: i \leq degree f

and hnorm: *normal-poly* f

shows 0 < coeff f i

<proof>

```

lemma normal-divide-x:
  fixes f::real poly
  assumes hnorm: normal-poly (f*[:0,1:])
  shows normal-poly f
  ⟨proof⟩

lemma normal-mult-x:
  fixes f::real poly
  assumes hnorm: normal-poly f
  shows normal-poly (f * [:0, 1:])
  ⟨proof⟩

lemma normal-poly-general-coeff-mult:
  fixes f::real poly
  assumes normal-poly f and  $h \leq j$ 
  shows  $\text{coeff } f (h+1) * \text{coeff } f (j+1) \geq \text{coeff } f h * \text{coeff } f (j+2)$ 
  ⟨proof⟩

lemma normal-mult:
  fixes f g::real poly
  assumes hf: normal-poly f and hg: normal-poly g
  defines df  $\equiv$  degree f and dg  $\equiv$  degree g
  shows normal-poly (f*g)
  ⟨proof⟩

lemma normal-poly-of-roots:
  fixes p::real poly
  assumes  $\bigwedge z. \text{poly } (\text{map-poly } \text{complex-of-real } p) z = 0$ 
     $\implies \text{Re } z \leq 0 \wedge 4 * (\text{Re } z)^2 \geq (\text{cmod } z)^2$ 
    and lead-coeff p = 1
  shows normal-poly p
  ⟨proof⟩

lemma normal-changes:
  fixes f::real poly
  assumes hf: normal-poly f and hx:  $x > 0$ 
  defines df  $\equiv$  degree f
  shows changes (coeffs (f*[:-x,1:])) = 1
  ⟨proof⟩

```

end

5 Proof of the theorem of three circles

```

theory Three-Circles
  imports Bernstein Normal-Poly
begin

```


The theorem of three circles is a result in real algebraic geometry about the number of real roots in an interval. It says if the number of roots in certain circles in the complex plane are zero or one then the number of roots in the circles is equal to the sign changes of the Bernstein coefficients on that interval for which the circles intersect the real line. This can then be used to determine if an interval has a real root in the bisection procedure, which is more efficient than Descartes' rule of signs.

The proof here follows Theorem 10.50 in Basu, S., Pollack, R., Roy, M.-F.: Algorithms in Real Algebraic Geometry. Springer Berlin Heidelberg, Berlin, Heidelberg (2016).

This theorem has also been formalised in Coq [4]. The relationship between this theorem and root isolation has been elaborated in Eigenwillig's PhD thesis [2].

5.1 No sign changes case

declare *degree-pcompose*[*simp del*]

corollary *descartes-sign-zero*:

fixes *p::real poly*
assumes $\bigwedge x::\text{complex. } \text{poly}(\text{map-poly of-real } p) x = 0 \implies \text{Re } x \leq 0$
and *lead-coeff* *p* = 1
shows *coeff* *p* *i* ≥ 0
<proof>

definition *circle-01-diam* :: *complex set where*

circle-01-diam =
 $\{x. \text{cmod}(x - (\text{of-nat } 1 :: \text{complex}) / (\text{of-nat } 2)) < (\text{real } 1) / (\text{real } 2)\}$

lemma *pos-real-map*:

$\{x::\text{complex. } 1 / x \in (\lambda x. x + 1) \text{ ' } \{x. 0 < \text{Re } x\}\} = \text{circle-01-diam}$
<proof>

lemma *one-circle-01*: **fixes** *P::real poly* **assumes** *hP*: *degree* *P* $\leq p$ **and** *P* $\neq 0$

and *proots-count* (*map-poly of-real* *P*) *circle-01-diam* = 0
shows *Bernstein-changes-01* *p* *P* = 0
<proof>

definition *circle-diam* :: *real \Rightarrow real \Rightarrow complex set where*

circle-diam *l r* = $\{x. \text{cmod}((x - l) - (r - l) / 2) < (r - l) / 2\}$

lemma *circle-diam-rescale*: **assumes** *l* < *r*

shows *circle-diam* *l r* = $(\lambda x. (x * (r - l) + l)) \text{ ' } \text{circle-01-diam}$
<proof>

lemma *one-circle*: **fixes** *P::real poly* **assumes** *l* < *r*

and *proots-count* (*map-poly of-real* *P*) (*circle-diam* *l r*) = 0

and $P \neq 0$
and $\text{degree } P \leq p$
shows *Bernstein-changes* $p \ l \ r \ P = 0$
 ⟨*proof*⟩

5.2 One sign change case

definition *upper-circle-01* :: *complex set* **where**
 $\text{upper-circle-01} = \{x. \text{cmod } (x - (1/2 + \text{sqrt}(3)/6 * i)) < \text{sqrt } 3 / 3\}$

lemma *upper-circle-map*:
 $\{x::\text{complex}. 1 / x \in (\lambda x. x + 1) \text{ ' } \{x. \text{Im } x < \text{sqrt } 3 * \text{Re } x\}\} = \text{upper-circle-01}$
 ⟨*proof*⟩

definition *lower-circle-01* :: *complex set* **where**
 $\text{lower-circle-01} = \{x. \text{cmod } (x - (1/2 - \text{sqrt}(3)/6 * i)) < \text{sqrt } 3 / 3\}$

lemma *cnj-upper-circle-01*: $\text{cnj ' upper-circle-01} = \text{lower-circle-01}$
 ⟨*proof*⟩

lemma *lower-circle-map*:
 $\{x::\text{complex}. 1 / x \in (\lambda x. x + 1) \text{ ' } \{x. \text{Im } x > -\text{sqrt } 3 * \text{Re } x\}\} = \text{lower-circle-01}$
 ⟨*proof*⟩

lemma *two-circles-01*:
fixes $P::\text{real poly}$
assumes $hP: \text{degree } P \leq p$ **and** $hP0: P \neq 0$ **and** $hp0: p \neq 0$
and $h: \text{roots-count } (\text{map-poly of-real } P)$
 $(\text{upper-circle-01} \cup \text{lower-circle-01}) = 1$
shows *Bernstein-changes-01* $p \ P = 1$
 ⟨*proof*⟩

definition *upper-circle* :: *real* \Rightarrow *real* \Rightarrow *complex set* **where**
 $\text{upper-circle } l \ r = \{x::\text{complex}.$
 $\text{cmod } ((x - \text{of-real } l) / (\text{of-real } (r - l)) - (1/2 + \text{of-real } (\text{sqrt}(3))/6 * i)) < \text{sqrt } 3 / 3\}$

lemma *upper-circle-rescale*: **assumes** $l < r$
shows $\text{upper-circle } l \ r = (\lambda x. (x * (r - l) + l)) \text{ ' } \text{upper-circle-01}$
 ⟨*proof*⟩

definition *lower-circle* :: *real* \Rightarrow *real* \Rightarrow *complex set* **where**
 $\text{lower-circle } l \ r = \{x::\text{complex}.$
 $\text{cmod } ((x - \text{of-real } l) / (\text{of-real } (r - l)) - (1/2 - \text{of-real } (\text{sqrt}(3))/6 * i)) < \text{sqrt } 3 / 3\}$

lemma *lower-circle-rescale*:
assumes $l < r$
shows $\text{lower-circle } l \ r = (\lambda x. (x * (r - l) + l)) \text{ ' } \text{lower-circle-01}$

<proof>

lemma *two-circles*:

fixes $P::\text{real poly}$ **and** $l\ r::\text{real}$

assumes $hlr: l < r$

and $hP: \text{degree } P \leq p$

and $hP0: P \neq 0$

and $hp0: p \neq 0$

and $h: \text{roots-count } (\text{map-poly of-real } P)$

$(\text{upper-circle } l\ r \cup \text{lower-circle } l\ r) = 1$

shows $\text{Bernstein-changes } p\ l\ r\ P = 1$

<proof>

5.3 The theorem of three circles

theorem *three-circles*:

fixes $P::\text{real poly}$ **and** $l\ r::\text{real}$

assumes $l < r$

and $hP: \text{degree } P \leq p$

and $hP0: P \neq 0$

and $hp0: p \neq 0$

shows $\text{roots-count } (\text{map-poly of-real } P) (\text{circle-diam } l\ r) = 0 \implies$

$\text{Bernstein-changes } p\ l\ r\ P = 0$

and $\text{roots-count } (\text{map-poly of-real } P)$

$(\text{upper-circle } l\ r \cup \text{lower-circle } l\ r) = 1 \implies$

$\text{Bernstein-changes } p\ l\ r\ P = 1$

<proof>

end

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