

Stirling's formula

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Abstract

This work contains a proof of Stirling's formula both for the factorial $n! \sim \sqrt{2\pi n} (n/e)^n$ on natural numbers and the real Gamma function $\Gamma(x) \sim \sqrt{2\pi/x} (x/e)^x$. The proof is based on work by Graham Jameson [3].

This is then extended to the full asymptotic expansion

$$\log \Gamma(z) = \left(z - \frac{1}{2}\right) \log z - z + \frac{1}{2} \log(2\pi) + \sum_{k=1}^{n-1} \frac{B_{k+1}}{k(k+1)} z^{-k} - \frac{1}{n} \int_0^\infty B_n([t])(t+z)^{-n} dt$$

uniformly for all complex $z \neq 0$ in the cone $\arg(z) \leq \alpha$ for any $\alpha \in (0, \pi)$, with which the above asymptotic relation for Γ is also extended to complex arguments.

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1 Stirling's Formula

```
theory Stirling-Formula
imports
  HOL-Analysis.Analysis
  Landau-Symbols.Landau-More
begin
```

context
begin

First, we define the S_n^* from Jameson's article:

private definition $S' :: \text{nat} \Rightarrow \text{real} \Rightarrow \text{real}$ **where**
 $S' n x = 1/(2*x) + (\sum r=1..<n. 1/(of-nat r+x)) + 1/(2*(n+x))$

Next, the trapezium (also called T in Jameson's article):

private definition $T :: \text{real} \Rightarrow \text{real}$ **where**
 $T x = 1/(2*x) + 1/(2*(x+1))$

Now we define The difference $\Delta(x)$:

private definition $D :: \text{real} \Rightarrow \text{real}$ **where**
 $D x = T x - \ln(x+1) + \ln x$

private lemma S' -telescope-trapezium:

assumes $n > 0$
shows $S' n x = (\sum r<n. T (of-nat r+x))$
proof (*cases n*)
case (*Suc m*)
hence $m: \text{Suc } m = n$ **by** *simp*
have $(\sum r<n. T (of-nat r+x)) =$
 $(\sum r<\text{Suc } m. 1 / (2 * \text{real } r + 2 * x)) + (\sum r<n. 1 / (2 * \text{real } (\text{Suc } r)$
 $+ 2 * x))$
unfolding m **by** (*simp add: T-def sum.distrib algebra-simps*)
also have $(\sum r<\text{Suc } m. 1 / (2 * \text{real } r + 2 * x)) =$
 $1/(2*x) + (\sum r<m. 1 / (2 * \text{real } (\text{Suc } r) + 2 * x))$ (**is - = ?a + ?S**)
by (*subst sum.lessThan-Suc-shift*) *simp*
also have $(\sum r<n. 1 / (2 * \text{real } (\text{Suc } r) + 2 * x)) =$
 $?S + 1 / (2*(\text{real } m + x + 1))$ (**is - = - + ?b**) **by** (*simp add: Suc*)
also have $?a + ?S + (?S + ?b) = 2*?S + ?a + ?b$ **by** (*simp add: add-ac*)
also have $2 * ?S = (\sum r=0..<m. 1 / (\text{real } (\text{Suc } r) + x))$
unfolding *sum-distrib-left* **by** (*intro sum.cong*) (*auto simp add: divide-simps*)
also have $(\sum r=0..<m. 1 / (\text{real } (\text{Suc } r) + x)) = (\sum r=\text{Suc } 0..<\text{Suc } m. 1 /$
 $(\text{real } r + x))$
by (*subst sum.atLeast-Suc-lessThan-Suc-shift*) *simp-all*
also have $\dots = (\sum r=1..<n. 1 / (\text{real } r + x))$ **unfolding** m **by** *simp*
also have $\dots + ?a + ?b = S' n x$ **by** (*simp add: S'-def Suc*)
finally show *?thesis ..*
qed (*insert assms, simp-all*)

private lemma *stirling-trapezium*:

assumes $x: (x::\text{real}) > 0$
shows $D x \in \{0 .. 1/(12*x^2) - 1/(12 * (x+1)^2)\}$
proof -
define y **where** $y = 1 / (2*x + 1)$
from x **have** $y > 0$ $y < 1$ **by** (*simp-all add: divide-simps y-def*)

from x **have** $D x = T x - \ln ((x + 1) / x)$ **by** (*subst ln-div*) (*simp-all add: D-def*)
also from x **have** $(x + 1) / x = 1 + 1 / x$ **by** (*simp add: field-simps*)
finally have $D: D x = T x - \ln (1 + 1/x)$.

from y **have** $(\lambda n. y * y^{\wedge} n) \text{ sums } (y * (1 / (1 - y)))$
by (*intro geometric-sums sums-mult*) *simp-all*
hence $(\lambda n. y^{\wedge} \text{Suc } n) \text{ sums } (y / (1 - y))$ **by** *simp*
also from x **have** $y / (1 - y) = 1 / (2*x)$ **by** (*simp add: y-def divide-simps*)
finally have $*$: $(\lambda n. y^{\wedge} \text{Suc } n) \text{ sums } (1 / (2*x))$.

from y **have** $(\lambda n. (-y) * (-y)^{\wedge} n) \text{ sums } ((-y) * (1 / (1 - (-y))))$
by (*intro geometric-sums sums-mult*) *simp-all*
hence $(\lambda n. (-y)^{\wedge} \text{Suc } n) \text{ sums } (-y / (1 + y))$ **by** *simp*
also from x **have** $y / (1 + y) = 1 / (2*(x+1))$ **by** (*simp add: y-def divide-simps*)
finally have $**$: $(\lambda n. (-y)^{\wedge} \text{Suc } n) \text{ sums } (-1 / (2*(x+1)))$.

from *sums-diff*[*OF * ***] **have** $\text{sum1}: (\lambda n. y^{\wedge} \text{Suc } n - (-y)^{\wedge} \text{Suc } n) \text{ sums } T x$
by (*simp add: T-def*)

from y **have** $\text{abs } y < 1$ $\text{abs } (-y) < 1$ **by** *simp-all*
from *sums-diff*[*OF this* [*THEN ln-series*]]
have $(\lambda n. y^{\wedge} n / \text{real } n - (-y)^{\wedge} n / \text{real } n) \text{ sums } (\ln (1 + y) - \ln (1 - y))$
by *simp*
also from y **have** $\ln (1 + y) - \ln (1 - y) = \ln ((1 + y) / (1 - y))$ **by** (*simp add: ln-div*)
also from x **have** $(1 + y) / (1 - y) = 1 + 1/x$ **by** (*simp add: divide-simps y-def*)
finally have $(\lambda n. y^{\wedge} n / \text{real } n - (-y)^{\wedge} n / \text{real } n) \text{ sums } \ln (1 + 1/x)$.
hence $\text{sum2}: (\lambda n. y^{\wedge} \text{Suc } n / \text{real } (\text{Suc } n) - (-y)^{\wedge} \text{Suc } n / \text{real } (\text{Suc } n)) \text{ sums } \ln (1 + 1/x)$
by (*subst sums-Suc-iff*) *simp*

have $\ln (1 + 1/x) \leq T x$
proof (*rule sums-le* [*OF - sum2 sum1*])
fix $n :: \text{nat}$
show $y^{\wedge} \text{Suc } n / \text{real } (\text{Suc } n) - (-y)^{\wedge} \text{Suc } n / \text{real } (\text{Suc } n) \leq y^{\wedge} \text{Suc } n - (-y)^{\wedge} \text{Suc } n$
proof (*cases even n*)
case *True*
hence $\text{eq}: A - (-y)^{\wedge} \text{Suc } n / B = A + y^{\wedge} \text{Suc } n / B$ $A - (-y)^{\wedge} \text{Suc } n = A + y^{\wedge} \text{Suc } n$
for $A B$ **by** *simp-all*
from y **show** *?thesis* **unfolding** *eq*
by (*intro add-mono*) (*auto simp: divide-simps*)
qed *simp-all*
qed
hence $D x \geq 0$ **by** (*simp add: D*)

define c **where** $c = (\lambda n. \text{if even } n \text{ then } 2 * (1 - 1 / \text{real } (\text{Suc } n)) \text{ else } 0)$
note $\text{sums-diff}[OF \text{ sum1 sum2}]$
also have $(\lambda n. y^{\wedge} \text{Suc } n - (-y)^{\wedge} \text{Suc } n - (y^{\wedge} \text{Suc } n / \text{real } (\text{Suc } n) - (-y)^{\wedge} \text{Suc } n / \text{real } (\text{Suc } n))) = (\lambda n. c \ n * y^{\wedge} \text{Suc } n)$
by $(\text{intro ext}) (\text{simp add: c-def algebra-simps})$
finally have $\text{sum3}: (\lambda n. c \ n * y^{\wedge} \text{Suc } n) \text{ sums } D \ x$ **by** $(\text{simp add: } D)$

from y **have** $(\lambda n. y^{\wedge} 2 * (\text{of-nat } (\text{Suc } n) * y^{\wedge} n)) \text{ sums } (y^{\wedge} 2 * (1 / (1 - y)^{\wedge} 2))$
by $(\text{intro sums-mult geometric-deriv-sums}) \text{ simp-all}$
hence $(\lambda n. \text{of-nat } (\text{Suc } n) * y^{\wedge} (n+2)) \text{ sums } (y^{\wedge} 2 / (1 - y)^{\wedge} 2)$
by $(\text{simp add: algebra-simps power2-eq-square})$
also from x **have** $y^{\wedge} 2 / (1 - y)^{\wedge} 2 = 1 / (4 * x^{\wedge} 2)$ **by** $(\text{simp add: y-def divide-simps})$
finally have $*$: $(\lambda n. \text{real } (\text{Suc } n) * y^{\wedge} (\text{Suc } (\text{Suc } n))) \text{ sums } (1 / (4 * x^2))$ **by** simp

from y **have** $(\lambda n. y^{\wedge} 2 * (\text{of-nat } (\text{Suc } n) * (-y)^{\wedge} n)) \text{ sums } (y^{\wedge} 2 * (1 / (1 - (-y)^{\wedge} 2)))$
by $(\text{intro sums-mult geometric-deriv-sums}) \text{ simp-all}$
hence $(\lambda n. \text{of-nat } (\text{Suc } n) * (-y)^{\wedge} (n+2)) \text{ sums } (y^{\wedge} 2 / (1 + y)^{\wedge} 2)$
by $(\text{simp add: algebra-simps power2-eq-square})$
also from x **have** $y^{\wedge} 2 / (1 + y)^{\wedge} 2 = 1 / (2^{\wedge} 2 * (x+1)^{\wedge} 2)$
unfolding $\text{power-mult-distrib [symmetric]}$ **by** $(\text{simp add: y-def divide-simps add-ac})$
finally have $**$: $(\lambda n. \text{real } (\text{Suc } n) * (-y)^{\wedge} (\text{Suc } (\text{Suc } n))) \text{ sums } (1 / (4 * (x + 1)^2))$ **by** simp

define d **where** $d = (\lambda n. \text{if even } n \text{ then } 2 * \text{real } n \text{ else } 0)$
note $\text{sums-diff}[OF * **]$
also have $(\lambda n. \text{real } (\text{Suc } n) * y^{\wedge} (\text{Suc } (\text{Suc } n)) - \text{real } (\text{Suc } n) * (-y)^{\wedge} (\text{Suc } (\text{Suc } n))) = (\lambda n. d (\text{Suc } n) * y^{\wedge} \text{Suc } (\text{Suc } n))$
by $(\text{intro ext}) (\text{simp-all add: d-def})$
finally have $(\lambda n. d \ n * y^{\wedge} \text{Suc } n) \text{ sums } (1 / (4 * x^2) - 1 / (4 * (x + 1)^2))$
by $(\text{subst (asm) sums-Suc-iff}) (\text{simp add: d-def})$
from $\text{sums-mult}[OF \text{ this, of } 1/3] \ x$
have $\text{sum4}: (\lambda n. d \ n / 3 * y^{\wedge} \text{Suc } n) \text{ sums } (1 / (12 * x^{\wedge} 2) - 1 / (12 * (x + 1)^{\wedge} 2))$
by $(\text{simp add: field-simps})$

have $D \ x \leq (1 / (12 * x^{\wedge} 2) - 1 / (12 * (x + 1)^{\wedge} 2))$
proof $(\text{intro sums-le [OF - sum3 sum4] allI})$
fix $n :: \text{nat}$
define $c' :: \text{nat} \Rightarrow \text{real}$
where $c' = (\lambda n. \text{if odd } n \vee n = 0 \text{ then } 0 \text{ else if } n = 2 \text{ then } 4/3 \text{ else } 2)$
show $c \ n * y^{\wedge} \text{Suc } n \leq d \ n / 3 * y^{\wedge} \text{Suc } n$
proof $(\text{intro mult-right-mono})$
have $c \ n \leq c' \ n$ **by** $(\text{simp add: c-def c'-def})$
also consider $n = 0 \mid n = 1 \mid n = 2 \mid n \geq 3$ **by** force

hence $c' n \leq d n / 3$ **by cases** (*simp-all add: c'-def d-def*)
finally show $c n \leq d n / 3$.
qed (*insert y, simp*)
qed

with $\langle D x \geq 0 \rangle$ **show** *?thesis by simp*
qed

The following functions correspond to the $p_n(x)$ from the article. The special case $n = 0$ would not, strictly speaking, be necessary, but it allows some theorems to work even without the precondition $n \neq 0$:

private definition $p :: nat \Rightarrow real \Rightarrow real$ **where**
 $p n x = (if\ n = 0\ then\ 1/x\ else\ (\sum\ r < n.\ D\ (real\ r + x)))$

We can write the Digamma function in terms of S' :

private lemma *S'-LIMSEQ-Digamma*:

assumes $x: x \neq 0$
shows $(\lambda n. \ln (real\ n) - S' n x - 1/(2*x)) \longrightarrow Digamma\ x$
proof –
define c **where** $c = (\lambda n. \ln (real\ n) - (\sum\ r < n.\ inverse\ (x + real\ r)))$
have *eventually* $(\lambda n. 1 / (2 * (x + real\ n)) = c n - (\ln (real\ n) - S' n x - 1/(2*x)))$ *at-top*
using *eventually-gt-at-top[of 0::nat]*
proof *eventually-elim*
fix $n :: nat$
assume $n: n > 0$
have $c n - (\ln (real\ n) - S' n x - 1/(2*x)) =$
 $-(\sum\ r < n.\ inverse\ (real\ r + x)) + (1/x + (\sum\ r = 1..<n.\ inverse\ (real\ r + x))) + 1/(2*(real\ n + x))$
using x **by** (*simp add: S'-def c-def field-simps*)
also have $1/x + (\sum\ r = 1..<n.\ inverse\ (real\ r + x)) = (\sum\ r < n.\ inverse\ (real\ r + x))$
unfolding *lessThan-atLeast0* **using** n
by (*subst (2) sum.atLeast-Suc-lessThan*) (*simp-all add: field-simps*)
finally show $1 / (2 * (x + real\ n)) = c n - (\ln (real\ n) - S' n x - 1/(2*x))$
by *simp*
qed
moreover have $(\lambda n. 1 / (2 * (x + real\ n))) \longrightarrow 0$
by (*rule real-tendsto-divide-at-top tendsto-const filterlim-tendsto-pos-mult-at-top filterlim-tendsto-add-at-top filterlim-real-sequentially | simp*)
ultimately have $(\lambda n. c n - (\ln (real\ n) - S' n x - 1/(2*x))) \longrightarrow 0$
by (*blast intro: Lim-transform-eventually*)
from *tendsto-minus[OF this]* **have** $(\lambda n. (\ln (real\ n) - S' n x - 1/(2*x)) - c n) \longrightarrow 0$ **by** *simp*
moreover from *Digamma-LIMSEQ[OF x]* **have** $c \longrightarrow Digamma\ x$ **by** (*simp add: c-def*)
ultimately show $(\lambda n. \ln (real\ n) - S' n x - 1/(2*x)) \longrightarrow Digamma\ x$
by (*rule Lim-transform [rotated]*)
qed

Moreover, we can give an expansion of S' with the p as variation terms.

private lemma S' -approx:

$$S' n x = \ln (\text{real } n + x) - \ln x + p n x$$

proof (cases $n = 0$)

case $True$

thus ?thesis **by** (simp add: p-def S' -def)

next

case $False$

$$\text{hence } S' n x = (\sum r < n. T (\text{real } r + x))$$

by (subst S' -trapezium) simp-all

$$\text{also have } \dots = (\sum r < n. \ln (\text{real } r + x + 1) - \ln (\text{real } r + x) + D (\text{real } r + x))$$

by (simp add: D-def)

$$\text{also have } \dots = (\sum r < n. \ln (\text{real } (\text{Suc } r) + x) - \ln (\text{real } r + x)) + p n x$$

using $False$ **by** (simp add: sum.distrib add-ac p-def)

$$\text{also have } (\sum r < n. \ln (\text{real } (\text{Suc } r) + x) - \ln (\text{real } r + x)) = \ln (\text{real } n + x) - \ln x$$

by (subst sum-lessThan-telescope) simp-all

finally show ?thesis .

qed

We define the limit of the p (simply called $p(x)$ in Jameson's article):

private definition $P :: \text{real} \Rightarrow \text{real}$ **where**

$$P x = (\sum n. D (\text{real } n + x))$$

private lemma D -summable:

assumes $x: x > 0$

shows summable $(\lambda n. D (\text{real } n + x))$

proof –

$$\text{have } *: \text{summable } (\lambda n. 1 / (12 * (x + \text{real } n)^2) - 1 / (12 * (x + \text{real } (\text{Suc } n))^2))$$

by (rule telescope-summable' real-tendsto-divide-at-top tendsto-const
filterlim-tendsto-pos-mult-at-top filterlim-pow-at-top
filterlim-tendsto-add-at-top filterlim-real-sequentially | simp)+

show summable $(\lambda n. D (\text{real } n + x))$

proof (rule summable-comparison-test[OF *], rule exI[of - 2], safe)

fix $n :: \text{nat}$ **assume** $n \geq 2$

$$\text{show } \text{norm } (D (\text{real } n + x)) \leq 1 / (12 * (x + \text{real } n)^2) - 1 / (12 * (x + \text{real } (\text{Suc } n))^2)$$

using stirling-trapezium[of $\text{real } n + x$] x **by** (auto simp: algebra-simps)

qed

qed

private lemma p -LIMSEQ:

assumes $x: x > 0$

shows $(\lambda n. p n x) \longrightarrow P x$

proof (rule Lim-transform-eventually)

from D -summable[OF x] **have** $(\lambda n. D (\text{real } n + x))$ sums $P x$ **unfolding** P -def

by (simp add: sums-iff)

then show $(\lambda n. \sum r < n. D (\text{real } r + x)) \longrightarrow P x$ **by** (*simp add: sums-def*)
moreover from *eventually-gt-at-top*[of 1]
show *eventually* $(\lambda n. (\sum r < n. D (\text{real } r + x)) = p n x)$ *at-top*
by *eventually-elim* (*auto simp: p-def*)
qed

This gives us an expansion of the Digamma function:

lemma *Digamma-approx*:

assumes $x: (x :: \text{real}) > 0$

shows $\text{Digamma } x = \ln x - 1 / (2 * x) - P x$

proof –

have *eventually* $(\lambda n. \ln (\text{inverse } (1 + x / \text{real } n)) + \ln x - 1/(2*x) - p n x = \ln (\text{real } n) - S' n x - 1/(2*x))$ *at-top*

using *eventually-gt-at-top*[of 1::nat]

proof *eventually-elim*

fix $n :: \text{nat}$ **assume** $n: n > 1$

have $\ln (\text{real } n) - S' n x = \ln ((\text{real } n) / (\text{real } n + x)) + \ln x - p n x$

using *assms n unfolding S'-approx by (subst ln-div) (auto simp: algebra-simps)*

also from n **have** $\text{real } n / (\text{real } n + x) = \text{inverse } (1 + x / \text{real } n)$ **by** (*simp add: field-simps*)

finally show $\ln (\text{inverse } (1 + x / \text{real } n)) + \ln x - 1/(2*x) - p n x = \ln (\text{real } n) - S' n x - 1/(2*x)$ **by** *simp*

qed

moreover have $(\lambda n. \ln (\text{inverse } (1 + x / \text{real } n)) + \ln x - 1/(2*x) - p n x) \longrightarrow \ln (\text{inverse } (1 + 0)) + \ln x - 1/(2*x) - P x$

by (*rule tendsto-intros p-LIMSEQ x real-tendsto-divide-at-top filterlim-real-sequentially | simp*)+

hence $(\lambda n. \ln (\text{inverse } (1 + x / \text{real } n)) + \ln x - 1/(2*x) - p n x) \longrightarrow \ln x - 1/(2*x) - P x$ **by** *simp*

ultimately have $(\lambda n. \ln (\text{real } n) - S' n x - 1 / (2 * x)) \longrightarrow \ln x - 1/(2*x) - P x$

by (*blast intro: Lim-transform-eventually*)

moreover from x **have** $(\lambda n. \ln (\text{real } n) - S' n x - 1 / (2 * x)) \longrightarrow \text{Digamma } x$

by (*intro S'-LIMSEQ-Digamma simp-all*)

ultimately show $\text{Digamma } x = \ln x - 1 / (2 * x) - P x$

by (*rule LIMSEQ-unique [rotated]*)

qed

Next, we derive some bounds on P :

private lemma *p-ge-0*: $x > 0 \implies p n x \geq 0$

using *stirling-trapezium*[of $\text{real } n + x$ **for** n]

by (*auto simp add: p-def intro!: sum-nonneg*)

private lemma *P-ge-0*: $x > 0 \implies P x \geq 0$

by (*rule tendsto-lowerbound[OF p-LIMSEQ]*)

(*insert p-ge-0*[of x], *simp-all*)

private lemma *p-upper-bound*:

assumes $x > 0$ $n > 0$

shows $p\ n\ x \leq 1/(12*x^2)$

proof –

from *assms* **have** $p\ n\ x = (\sum r < n. D\ (real\ r + x))$

by (*simp add: p-def*)

also have $\dots \leq (\sum r < n. 1/(12*(real\ r + x)^2) - 1/(12 * (real\ (Suc\ r) + x)^2))$

using *stirling-trapezium[of real r + x for r] assms*

by (*intro sum-mono*) (*simp add: add-ac*)

also have $\dots = 1 / (12 * x^2) - 1 / (12 * (real\ n + x)^2)$

by (*subst sum-lessThan-telescope'*) *simp*

also have $\dots \leq 1 / (12 * x^2)$ **by** *simp*

finally show *?thesis* .

qed

private lemma *P-upper-bound*:

assumes $x > 0$

shows $P\ x \leq 1/(12*x^2)$

proof (*rule tendsto-upperbound*)

show *eventually* $(\lambda n. p\ n\ x \leq 1 / (12 * x^2))$ *at-top*

using *eventually-gt-at-top[of 0]*

by *eventually-elim* (*use p-upper-bound[of x] assms in auto*)

show $(\lambda n. p\ n\ x) \longrightarrow P\ x$

by (*simp add: assms p-LIMSEQ*)

qed *auto*

We can now use this approximation of the Digamma function to approximate *ln-Gamma* (since the former is the derivative of the latter). We therefore define the function *g* from Jameson's article, which measures the difference between *ln-Gamma* and its approximation:

private definition $g :: real \Rightarrow real$ **where**

$g\ x = ln-Gamma\ x - (x - 1/2) * ln\ x + x$

private lemma *DERIV-g*: $x > 0 \implies (g\ \text{has-field-derivative}\ -P\ x)\ (at\ x)$

unfolding *g-def* [*abs-def*] **using** *Digamma-approx*[*of x*]

by (*auto intro!: derivative-eq-intros simp: field-simps*)

private lemma *isCont-P*:

assumes $x > 0$

shows *isCont* $P\ x$

proof –

define $D' :: real \Rightarrow real$

where $D' = (\lambda x. - 1 / (2 * x^2 * (x+1)^2))$

have *DERIV-D*: $(D\ \text{has-field-derivative}\ D'\ x)\ (at\ x)$ **if** $x > 0$ **for** x

unfolding *D-def* [*abs-def*] *D'-def* *T-def*

by (*insert that, (rule derivative-eq-intros refl | simp)+*)

(*simp add: power2-eq-square divide-simps, (simp add: algebra-simps)*?)

note *this* [*THEN DERIV-chain2, derivative-intros*]


```

have (P has-field-derivative ( $\sum n. D' (real\ n + x)$ )) (at x)
  unfolding P-def [abs-def]
proof (rule has-field-derivative-series')
  show convex {x/2<..} by simp
next
  fix n :: nat and y :: real assume y: y  $\in$  {x/2<..}
  with assms have y > 0 by simp
  thus (( $\lambda a. D (real\ n + a)$ ) has-real-derivative  $D' (real\ n + y)$ ) (at y within
{x/2<..})
    by (auto intro!: derivative-eq-intros)
next
  from assms D-summable[of x] show summable ( $\lambda n. D (real\ n + x)$ ) by simp
next
  show uniformly-convergent-on {x/2<..} ( $\lambda n\ x. \sum i < n. D' (real\ i + x)$ )
  proof (rule Weierstrass-m-test')
    fix n :: nat and y :: real
    assume y: y  $\in$  {x/2<..}
    with assms have y > 0 by auto
    have norm ( $D' (real\ n + y)$ ) = (1 / (2 * (y + real n)2)) * (1 / (y + real
(Suc n))2)
      by (simp add: D'-def add-ac)
    also from y assms have ...  $\leq$  (1 / (2 * (x/2)2)) * (1 / (real (Suc n))2)
      by (intro mult-mono divide-left-mono power-mono) simp-all
    also have 1 / (real (Suc n))2 = inverse ((real (Suc n))2) by (simp add:
field-simps)
    finally show norm ( $D' (real\ n + y)$ )  $\leq$  (1 / (2 * (x/2)2)) * inverse ((real
(Suc n))2) .
  next
    show summable ( $\lambda n. (1 / (2 * (x/2)2)) * inverse ((real (Suc n))2)$ )
      by (subst summable-Suc-iff, intro summable-mult inverse-power-summable)
  simp-all
  qed
  qed (insert assms, simp-all add: interior-open)
  thus ?thesis by (rule DERIV-isCont)
qed

```

```

private lemma P-continuous-on [THEN continuous-on-subset]: continuous-on {0<..}
P
  by (intro continuous-at-imp-continuous-on ballI isCont-P) auto

```

```

private lemma P-integrable:
  assumes a: a > 0
  shows P integrable-on {a..}
proof -
  define f where f = ( $\lambda n\ x. \text{if } x \in \{a..real\ n\} \text{ then } P\ x \text{ else } 0$ )
  show P integrable-on {a..}
  proof (rule dominated-convergence)
    fix n :: nat

```

```

from  $a$  have  $P$  integrable-on  $\{a..real\ n\}$ 
  by (intro integrable-continuous-real P-continuous-on) auto
hence  $f\ n$  integrable-on  $\{a..real\ n\}$ 
  by (rule integrable-eq) (simp add: f-def)
thus  $f\ n$  integrable-on  $\{a..\}$ 
  by (rule integrable-on-superset) (auto simp: f-def)
next
  fix  $n :: nat$ 
  show  $norm\ (f\ n\ x) \leq of\ real\ (1/12) * (1 / x^2)$  if  $x \in \{a..\}$  for  $x$ 
    using  $a$  P-ge-0 P-upper-bound by (auto simp: f-def)
next
  show  $(\lambda x::real. of\ real\ (1/12) * (1 / x^2))$  integrable-on  $\{a..\}$ 
    using has-integral-inverse-power-to-inf[of 2 a]  $a$ 
    by (intro integrable-on-cmult-left) auto
next
  show  $(\lambda n. f\ n\ x) \longrightarrow P\ x$  if  $x \in \{a..\}$  for  $x$ 
  proof -
    have eventually  $(\lambda n. real\ n \geq x)$  at-top
      using filterlim-real-sequentially by (simp add: filterlim-at-top)
    with that have eventually  $(\lambda n. f\ n\ x = P\ x)$  at-top
      by (auto elim!: eventually-mono simp: f-def)
    thus  $(\lambda n. f\ n\ x) \longrightarrow P\ x$  by (simp add: tendsto-eventually)
  qed
qed
qed

```

private definition $c :: real$ **where** $c = integral\ \{1..\}\ (\lambda x. -P\ x) + g\ 1$

We can now give bounds on g :

private lemma *g-bounds*:

```

assumes  $x: x \geq 1$ 
shows  $g\ x \in \{c..c + 1/(12*x)\}$ 
proof -
  from assms have int-nonneg: integral  $\{x..\}$   $P \geq 0$ 
    by (intro Henstock-Kurzweil-Integration.integral-nonneg P-integrable)
    (auto simp: P-ge-0)
  have int-upper-bound: integral  $\{x..\}$   $P \leq 1/(12*x)$ 
  proof (rule has-integral-le)
    from  $x$  show  $(P\ has\ integral\ integral\ \{x..\}\ P)\ \{x..\}$ 
      by (intro integrable-integral P-integrable) simp-all
    from  $x$  has-integral-mult-right[OF has-integral-inverse-power-to-inf[of 2 x], of 1/12]
      show  $((\lambda x. 1/(12*x^2))\ has\ integral\ (1/(12*x)))\ \{x..\}$  by (simp add: field-simps)
  qed (insert P-upper-bound x, simp-all)

note DERIV-g [THEN DERIV-chain2, derivative-intros]
from assms have int1: ((\lambda x. -P x) has-integral (g x - g 1))  $\{1..x\}$ 
  by (intro fundamental-theorem-of-calculus)

```

(auto simp: has-field-derivative-iff-has-vector-derivative [symmetric]
 intro!: derivative-eq-intros)
from x **have** $\text{int2: } ((\lambda x. -P x) \text{ has-integral integral } \{x..\} (\lambda x. -P x)) \{x..\}$
by (intro integrable-integral integrable-neg P-integrable) simp-all
from $\text{has-integral-Un[OF int1 int2]} x$
have $((\lambda x. -P x) \text{ has-integral } g x - g 1 - \text{integral } \{x..\} P) (\{1..x\} \cup \{x..\})$
by (simp add: max-def)
also from x **have** $\{1..x\} \cup \{x..\} = \{1..\}$ **by** auto
finally have $((\lambda x. -P x) \text{ has-integral } g x - g 1 - \text{integral } \{x..\} P) \{1..\}$.
moreover have $((\lambda x. -P x) \text{ has-integral integral } \{1..\} (\lambda x. -P x)) \{1..\}$
by (intro integrable-integral integrable-neg P-integrable) simp-all
ultimately have $g x - g 1 - \text{integral } \{x..\} P = \text{integral } \{1..\} (\lambda x. -P x)$
by (simp add: has-integral-unique)
hence $g x = c + \text{integral } \{x..\} P$ **by** (simp add: c-def algebra-simps)
with $\text{int-upper-bound int-nonneg}$ **show** $g x \in \{c..c + 1/(12*x)\}$ **by** simp
qed

Finally, we have bounds on $\ln\text{-Gamma}$, Gamma , and fact .

private lemma $\text{ln-Gamma-bounds-aux}$:

$x \geq 1 \implies \text{ln-Gamma } x \geq c + (x - 1/2) * \ln x - x$
 $x \geq 1 \implies \text{ln-Gamma } x \leq c + (x - 1/2) * \ln x - x + 1/(12*x)$
using $g\text{-bounds[of } x]$ **by** (simp-all add: g-def)

private lemma Gamma-bounds-aux :

assumes $x: x \geq 1$
shows $\text{Gamma } x \geq \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x$
 $\text{Gamma } x \leq \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x * \text{exp } (1/(12*x))$

proof –

have $\text{exp } (\text{ln-Gamma } x) \geq \text{exp } (c + (x - 1/2) * \ln x - x)$
by (subst exp-le-cancel-iff, rule ln-Gamma-bounds-aux) (simp add: x)
with x **show** $\text{Gamma } x \geq \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x$
by (simp add: Gamma-real-pos-exp exp-add exp-diff powr-def del: exp-le-cancel-iff)

next

have $\text{exp } (\text{ln-Gamma } x) \leq \text{exp } (c + (x - 1/2) * \ln x - x + 1/(12*x))$
by (subst exp-le-cancel-iff, rule ln-Gamma-bounds-aux) (simp add: x)
with x **show** $\text{Gamma } x \leq \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x * \text{exp } (1/(12*x))$
by (simp add: Gamma-real-pos-exp exp-add exp-diff powr-def del: exp-le-cancel-iff)

qed

private lemma $\text{Gamma-asympt-equiv-aux}$:

$\text{Gamma} \sim_{[\text{at-top}]} (\lambda x. \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x)$

proof (rule asympt-equiv-sandwich)

include $\text{asympt-equiv-notation}$

show $\text{eventually } (\lambda x. \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x \leq \text{Gamma } x) \text{ at-top}$
 $\text{eventually } (\lambda x. \text{exp } c * x \text{ powr } (x - 1/2) / \text{exp } x * \text{exp } (1/(12*x)) \geq \text{Gamma } x) \text{ at-top}$

using $\text{eventually-ge-at-top[of } 1::\text{real}]$

by (eventually-elim; use Gamma-bounds-aux in force)+
have $((\lambda x::\text{real}. \text{exp } (1 / (12 * x))) \longrightarrow \text{exp } 0) \text{ at-top}$

by (*rule tendsto-intros real-tendsto-divide-at-top filterlim-tendsto-pos-mult-at-top*) +
(simp-all add: filterlim-ident)
hence $(\lambda x. \exp (1 / (12 * x))) \sim (\lambda x. 1 :: \text{real})$
by (*intro asymp-equivI'*) *simp-all*
hence $(\lambda x. \exp c * x \text{ powr } (x - 1 / 2) / \exp x * 1) \sim$
 $(\lambda x. \exp c * x \text{ powr } (x - 1 / 2) / \exp x * \exp (1 / (12 * x)))$
by (*intro asymp-equiv-mult asymp-equiv-refl*) (*simp add: asymp-equiv-sym*)
thus $(\lambda x. \exp c * x \text{ powr } (x - 1 / 2) / \exp x) \sim$
 $(\lambda x. \exp c * x \text{ powr } (x - 1 / 2) / \exp x * \exp (1 / (12 * x)))$ **by** *simp*
qed *simp-all*

private lemma *exp-1-powr-real* [*simp*]: $\exp (1 :: \text{real}) \text{ powr } x = \exp x$
by (*simp add: powr-def*)

private lemma *fact-asymp-equiv-aux*:

fact \sim [*at-top*] $(\lambda x. \exp c * \text{sqrt} (\text{real } x) * (\text{real } x / \exp 1) \text{ powr } \text{real } x)$

proof –

include *asymp-equiv-notation*

have *fact* \sim $(\lambda n. \text{Gamma} (\text{real} (\text{Suc } n)))$ **by** (*simp add: Gamma-fact*)

also have *eventually* $(\lambda n. \text{Gamma} (\text{real} (\text{Suc } n)) = \text{real } n * \text{Gamma} (\text{real } n))$

at-top

using *eventually-gt-at-top* [*of 0 :: nat*]

by *eventually-elim* (*insert Gamma-plus1* [*of real n for n*],

auto simp: add-ac of-nat-in-nonpos-Ints-iff)

also have $(\lambda n. \text{Gamma} (\text{real } n)) \sim (\lambda n. \exp c * \text{real } n \text{ powr } (\text{real } n - 1 / 2) /$
 $\exp (\text{real } n))$

by (*rule asymp-equiv-compose'* [*OF Gamma-asymp-equiv-aux*] *filterlim-real-sequentially*) +

also have *eventually* $(\lambda n. \text{real } n * (\exp c * \text{real } n \text{ powr } (\text{real } n - 1 / 2) / \exp$
 $(\text{real } n)) =$

$\exp c * \text{sqrt} (\text{real } n) * (\text{real } n / \exp 1) \text{ powr } \text{real } n)$ *at-top*

using *eventually-gt-at-top* [*of 0 :: nat*]

proof *eventually-elim*

fix $n :: \text{nat}$ **assume** $n > 0$

thus $\text{real } n * (\exp c * \text{real } n \text{ powr } (\text{real } n - 1 / 2) / \exp (\text{real } n)) =$

$\exp c * \text{sqrt} (\text{real } n) * (\text{real } n / \exp 1) \text{ powr } \text{real } n$

by (*subst powr-diff*) (*simp-all add: powr-divide powr-half-sqrt field-simps*)

qed

finally show *?thesis* **by** – (*simp-all add: asymp-equiv-mult*)

qed

We still need to determine the constant term c , which we do using Wallis' product formula:

$$\prod_{n=1}^{\infty} \frac{4n^2}{4n^2 - 1} = \frac{\pi}{2}$$

private lemma *powr-mult-2*: $(x :: \text{real}) > 0 \implies x \text{ powr } (y * 2) = (x \wedge 2) \text{ powr } y$

by (*subst mult.commute, subst powr-powr* [*symmetric*]) (*simp add: powr-numeral*)

private lemma *exp-mult-2*: $\exp (y * 2 :: \text{real}) = \exp y * \exp y$

by (subst exp-add [symmetric]) simp

private lemma exp-c: exp c = sqrt (2*pi)

proof –

include asymp-equiv-notation

define p where p = (λn. ∏ k=1..n. (4*real k²) / (4*real k² - 1))

have p-0 [simp]: p 0 = 1 by (simp add: p-def)

have p-Suc: p (Suc n) = p n * (4 * real (Suc n)²) / (4 * real (Suc n)² - 1)

for n unfolding p-def by (subst prod.nat-ivl-Suc') simp-all

have p: p = (λn. 16ⁿ * fact n⁴ / (fact (2 * n))² / (2 * real n + 1))

proof

fix n :: nat

have p n = (∏ k=1..n. (2*real k)² / (2*real k - 1) / (2 * real k + 1))

unfolding p-def by (intro prod.cong refl) (simp add: field-simps power2-eq-square)

also have ... = (∏ k=1..n. (2*real k)² / (2*real k - 1)) / (∏ k=1..n. (2 * real (Suc k) - 1))

by (simp add: prod-dividef prod.distrib add-ac)

also have (∏ k=1..n. (2 * real (Suc k) - 1)) = (∏ k=Suc 1..Suc n. (2 * real k - 1))

by (subst prod.atLeast-Suc-atMost-Suc-shift) simp-all

also have ... = (∏ k=1..n. (2 * real k - 1)) * (2 * real n + 1)

by (induction n) (simp-all add: prod.nat-ivl-Suc')

also have (∏ k = 1..n. (2 * real k)² / (2 * real k - 1)) / ... =

(∏ k = 1..n. (2 * real k)² / (2 * real k - 1)²) / (2 * real n + 1)

unfolding power2-eq-square by (simp add: prod.distrib prod-dividef)

also have (∏ k = 1..n. (2 * real k)² / (2 * real k - 1)²) =

(∏ k = 1..n. (2 * real k)⁴ / ((2*real k)*(2 * real k - 1))²)

by (rule prod.cong) (simp-all add: power2-eq-square eval-nat-numeral)

also have ... = 16ⁿ * fact n⁴ / (∏ k=1..n. (2*real k) * (2*real k - 1))²

by (simp add: prod.distrib prod-dividef fact-prod prod-power-distrib [symmetric] prod-constant)

also have (∏ k=1..n. (2*real k) * (2*real k - 1)) = fact (2*n)

by (induction n) (simp-all add: algebra-simps prod.nat-ivl-Suc')

finally show p n = 16ⁿ * fact n⁴ / (fact (2 * n))² / (2 * real n + 1) .

qed

have p ~ (λn. 16ⁿ * fact n⁴ / (fact (2 * n))² / (2 * real n + 1))

by (simp add: p)

also have ... ~ (λn. 16ⁿ * (exp c * sqrt (real n) * (real n / exp 1) powr real n)⁴ /

(exp c * sqrt (real (2*n)) * (real (2*n) / exp 1) powr real (2*n))² /

(2 * real n + 1)) (is - ~ ?f)

by (intro asymp-equiv-mult asymp-equiv-divide asymp-equiv-refl mult-nat-left-at-top fact-asymp-equiv-aux asymp-equiv-power asymp-equiv-compose'[OF fact-asymp-equiv-aux])

simp-all

also have eventually (λn. ... n = exp c² / (4 + 2/n)) at-top

```

using eventually-gt-at-top[of 0::nat]
proof eventually-elim
  fix n :: nat assume n: n > 0
  have [simp]:  $16^{\hat{n}} = 4^{\hat{n}} * (4^{\hat{n}} :: real)$  by (simp add: power-mult-distrib
[symmetric])
  from n have ?f n =  $\exp c^{\hat{2}} * (n / (2*(2*n+1)))$ 
  by (simp add: power-mult-distrib divide-simps powr-mult real-sqrt-power-even)
  (simp add: field-simps power2-eq-square eval-nat-numeral power-mult-2
exp-mult-2 powr-realpow)
  also from n have ... =  $\exp c^{\hat{2}} / (4 + 2/n)$  by (simp add: field-simps)
  finally show ?f n = ... .
qed
also have  $(\lambda x. 4 + 2 / real\ x) \sim (\lambda x. 4)$ 
  by (subst asymp-equiv-add-right) auto
finally have  $p \longrightarrow \exp c^{\hat{2}} / 4$ 
  by (rule asymp-equivD-const) (simp-all add: asymp-equiv-divide)
moreover have  $p \longrightarrow \pi / 2$  unfolding p-def by (rule wallis)
ultimately have  $\exp c^{\hat{2}} / 4 = \pi / 2$  by (rule LIMSEQ-unique)
hence  $2 * \pi = \exp c^{\hat{2}}$  by simp
also have  $\sqrt{\exp c^{\hat{2}}} = \exp c$  by simp
finally show  $\exp c = \sqrt{2 * \pi}$  ..
qed

```

```

private lemma c:  $c = \ln (2*\pi) / 2$ 
proof -
  note exp-c [symmetric]
  also have  $\ln (\exp c) = c$  by simp
  finally show ?thesis by (simp add: ln-sqrt)
qed

```

This gives us the final bounds:

theorem Gamma-bounds:

```

assumes  $x \geq 1$ 
shows  $\Gamma x \geq \sqrt{2*\pi/x} * (x / \exp 1)^{x-1/2}$  (is ?th1)
   $\Gamma x \leq \sqrt{2*\pi/x} * (x / \exp 1)^{x-1/2} * \exp (1 / (12 * x))$  (is
?th2)

```

proof -

```

from assms have  $\exp c * x^{x-1/2} / \exp x = \sqrt{2*\pi/x} * (x / \exp 1)^{x-1/2}$ 
by (subst powr-diff)
  (simp add: exp-c real-sqrt-divide powr-divide powr-half-sqrt field-simps)
with Gamma-bounds-aux[OF assms] show ?th1 ?th2 by simp-all
qed

```

theorem ln-Gamma-bounds:

```

assumes  $x \geq 1$ 
shows  $\ln \Gamma x \geq \ln (2*\pi/x) / 2 + x * \ln x - x$  (is ?th1)
   $\ln \Gamma x \leq \ln (2*\pi/x) / 2 + x * \ln x - x + 1/(12*x)$  (is ?th2)
proof -

```

from *ln-Gamma-bounds-aux*[*OF* *assms*] *assms* **show** *?th1* *?th2*
by (*simp-all* *add: c field-simps ln-div*)
from *assms* **have** $\exp c * x \text{ powr } (x - 1/2) / \exp x = \text{sqrt } (2*\pi/x) * (x / \exp$
1) powr x
by (*subst powr-diff*)
(simp add: exp-c real-sqrt-divide powr-divide powr-half-sqrt field-simps)
qed

theorem *fact-bounds*:

assumes $n > 0$
shows $(\text{fact } n :: \text{real}) \geq \text{sqrt } (2*\pi*n) * (n / \exp 1) ^ n$ (**is** *?th1*)
 $(\text{fact } n :: \text{real}) \leq \text{sqrt } (2*\pi*n) * (n / \exp 1) ^ n * \exp (1 / (12 * n))$ (**is**
?th2)

proof –

from *assms* **have** $n: \text{real } n \geq 1$ **by** *simp*
from *assms* *Gamma-plus1*[*of* *real n*]
have $\text{real } n * \text{Gamma } (\text{real } n) = \text{Gamma } (\text{real } (\text{Suc } n))$
by (*simp add: of-nat-in-nonpos-Ints-iff add-ac*)
also **have** $\text{Gamma } (\text{real } (\text{Suc } n)) = \text{fact } n$ **by** (*subst Gamma-fact* [*symmetric*])
simp
finally **have** $*$: $\text{fact } n = \text{real } n * \text{Gamma } (\text{real } n)$ **by** *simp*

have $2*\pi/n = 2*\pi*n / n^2$ **by** (*simp add: power2-eq-square*)
also **have** $\text{sqrt } \dots = \text{sqrt } (2*\pi*n) / n$ **by** (*subst real-sqrt-divide*) *simp-all*
also **have** $\text{real } n * \dots = \text{sqrt } (2*\pi*n)$ **by** *simp*
finally **have** $**$: $\text{real } n * \text{sqrt } (2*\pi/\text{real } n) = \text{sqrt } (2*\pi*\text{real } n)$.

note $*$

also **note** *Gamma-bounds(2)*[*OF* *n*]
also **have** $\text{real } n * (\text{sqrt } (2 * \pi / \text{real } n) * (\text{real } n / \exp 1) \text{ powr } \text{real } n * \exp (1 / (12 * \text{real } n))) =$
 $(\text{real } n * \text{sqrt } (2*\pi/n)) * (n / \exp 1) \text{ powr } n * \exp (1 / (12 * n))$
by (*simp add: algebra-simps*)
also **from** *n* **have** $(\text{real } n / \exp 1) \text{ powr } \text{real } n = (\text{real } n / \exp 1) ^ n$
by (*subst powr-realpow*) *simp-all*
also **note** $**$
finally **show** *?th2* **by** – (*insert assms, simp-all*)

have $\text{sqrt } (2*\pi*n) * (n / \exp 1) \text{ powr } n = n * (\text{sqrt } (2*\pi/n) * (n / \exp 1) \text{ powr } n)$
by (*subst *** [*symmetric*]) (*simp add: field-simps*)
also **from** *assms* **have** $\dots \leq \text{real } n * \text{Gamma } (\text{real } n)$
by (*intro mult-left-mono Gamma-bounds(1)*) *simp-all*
also **from** *n* **have** $(\text{real } n / \exp 1) \text{ powr } \text{real } n = (\text{real } n / \exp 1) ^ n$
by (*subst powr-realpow*) *simp-all*
also **note** $*$ [*symmetric*]
finally **show** *?th1* .
qed

theorem *ln-fact-bounds*:
assumes $n > 0$
shows $\ln (\text{fact } n :: \text{real}) \geq \ln (2 * \pi * n) / 2 + n * \ln n - n$ (**is** ?th1)
 $\ln (\text{fact } n :: \text{real}) \leq \ln (2 * \pi * n) / 2 + n * \ln n - n + 1 / (12 * \text{real } n)$ (**is** ?th2)
proof –
have $\ln (\text{fact } n :: \text{real}) \geq \ln (\text{sqrt } (2 * \pi * \text{real } n) * (\text{real } n / \text{exp } 1) ^ n)$
using *fact-bounds(1)*[*OF* *assms*] *assms* **by** (*subst ln-le-cancel-iff*) *auto*
also from *assms* **have** $\ln (\text{sqrt } (2 * \pi * \text{real } n) * (\text{real } n / \text{exp } 1) ^ n) = \ln (2 * \pi * n) / 2 + n * \ln n - n$
by (*simp add: ln-mult ln-div ln-realpow ln-sqrt field-simps*)
finally show ?th1 .
next
have $\ln (\text{fact } n :: \text{real}) \leq \ln (\text{sqrt } (2 * \pi * \text{real } n) * (\text{real } n / \text{exp } 1) ^ n * \text{exp } (1 / (12 * \text{real } n)))$
using *fact-bounds(2)*[*OF* *assms*] *assms* **by** (*subst ln-le-cancel-iff*) *auto*
also from *assms* **have** $\dots = \ln (2 * \pi * n) / 2 + n * \ln n - n + 1 / (12 * \text{real } n)$
by (*simp add: ln-mult ln-div ln-realpow ln-sqrt field-simps*)
finally show ?th2 .
qed

theorem *Gamma-asymp-equiv*:
 $\Gamma x \sim_{[\text{at-top}]} (\lambda x. \text{sqrt } (2 * \pi / x) * (x / \text{exp } 1) ^ x :: \text{real})$
proof –
note *Gamma-asymp-equiv-aux*
also have *eventually* $(\lambda x. \text{exp } c * x ^ (x - 1 / 2) / \text{exp } x = \text{sqrt } (2 * \pi / x) * (x / \text{exp } 1) ^ x)$ *at-top*
using *eventually-gt-at-top*[*of 0::real*]
proof *eventually-elim*
fix $x :: \text{real}$ **assume** $x > 0$
thus $\text{exp } c * x ^ (x - 1 / 2) / \text{exp } x = \text{sqrt } (2 * \pi / x) * (x / \text{exp } 1) ^ x$
by (*subst powr-diff*)
(simp add: exp-c powr-half-sqrt powr-divide field-simps real-sqrt-divide)
qed
finally show ?thesis .
qed

theorem *fact-asymp-equiv*:
 $\text{fact } n \sim_{[\text{at-top}]} (\lambda n. \text{sqrt } (2 * \pi * n) * (n / \text{exp } 1) ^ n :: \text{real})$
proof –
note *fact-asymp-equiv-aux*
also have *eventually* $(\lambda n. \text{exp } c * \text{sqrt } (\text{real } n) = \text{sqrt } (2 * \pi * \text{real } n))$ *at-top*
using *eventually-gt-at-top*[*of 0::nat*] **by** *eventually-elim* (*simp add: exp-c real-sqrt-mult*)
also have *eventually* $(\lambda n. (n / \text{exp } 1) ^ n = (n / \text{exp } 1) ^ n)$ *at-top*
using *eventually-gt-at-top*[*of 0::nat*] **by** *eventually-elim* (*simp add: powr-realpow*)
finally show ?thesis .
qed

end

end

2 Complete asymptotics of the logarithmic Gamma function

theory *Gamma-Asymptotics*

imports

HOL-Complex-Analysis.Complex-Analysis

HOL-Real-Asymp.Real-Asymp

Bernoulli.Bernoulli-FPS

Bernoulli.Periodic-Bernpoly

Stirling-Formula

begin

2.1 Auxiliary Facts

lemma *stirling-limit-aux1*:

$((\lambda y. \text{Ln } (1 + z * \text{of-real } y) / \text{of-real } y) \longrightarrow z) \text{ (at-right } 0) \text{ for } z :: \text{complex}$

proof (*cases* $z = 0$)

case *True*

then show *?thesis* **by** *simp*

next

case *False*

have $((\lambda y. \ln (1 + z * \text{of-real } y)) \text{ has-vector-derivative } 1 * z) \text{ (at } 0)$

by (*rule* *has-vector-derivative-real-field*) (*auto intro!*: *derivative-eq-intros*)

then have $(\lambda y. (\text{Ln } (1 + z * \text{of-real } y) - \text{of-real } y * z) / \text{of-real } |y|) - 0 \rightarrow 0$

by (*auto simp add*: *has-vector-derivative-def has-derivative-def netlimit-at scaleR-conv-of-real field-simps*)

then have $((\lambda y. (\text{Ln } (1 + z * \text{of-real } y) - \text{of-real } y * z) / \text{of-real } |y|) \longrightarrow 0) \text{ (at-right } 0)$

by (*rule* *filterlim-mono[OF - - at-le]*) *simp-all*

also have *?this* $\longleftrightarrow ((\lambda y. \text{Ln } (1 + z * \text{of-real } y) / (\text{of-real } y) - z) \longrightarrow 0) \text{ (at-right } 0)$

using *eventually-at-right-less[of 0::real]*

by (*intro* *filterlim-cong refl*) (*auto elim!*: *eventually-mono simp: field-simps*)

finally show *?thesis* **by** (*simp only*: *LIM-zero-iff*)

qed

lemma *stirling-limit-aux2*:

$((\lambda y. y * \text{Ln } (1 + z / \text{of-real } y)) \longrightarrow z) \text{ at-top for } z :: \text{complex}$

using *stirling-limit-aux1* [*of z*] **by** (*subst* *filterlim-at-top-to-right*) (*simp add*: *field-simps*)

lemma *Union-atLeastAtMost*:

assumes $N > 0$

shows $(\bigcup n \in \{0..<N\}. \{\text{real } n.. \text{real } (n + 1)\}) = \{0.. \text{real } N\}$

proof (*intro* *equalityI subsetI*)

fix x **assume** $x \in \{0.. \text{real } N\}$

```

thus  $x \in (\bigcup n \in \{0..<N\}. \{real\ n..real\ (n + 1)\})$ 
proof (cases  $x = real\ N$ )
  case True
    with assms show ?thesis by (auto intro!: bexI[of - N - 1])
  next
    case False
    with  $x$  have  $x \geq 0\ x < real\ N$  by simp-all
    hence  $x \geq real\ (nat\ \lfloor x \rfloor)\ x \leq real\ (nat\ \lfloor x \rfloor + 1)$  by linarith+
    moreover from  $x$  have  $nat\ \lfloor x \rfloor < N$  by linarith
    ultimately have  $\exists n \in \{0..<N\}. x \in \{real\ n..real\ (n + 1)\}$ 
      by (intro bexI[of - nat\ \lfloor x \rfloor]) simp-all
    thus ?thesis by blast
  qed
qed auto

```

2.2 Cones in the complex plane

definition *complex-cone* :: $real \Rightarrow real \Rightarrow complex\ set$ **where**
complex-cone $a\ b = \{z. \exists y \in \{a..b\}. z = rcis\ (norm\ z)\ y\}$

abbreviation *complex-cone'* :: $real \Rightarrow complex\ set$ **where**
complex-cone' $a \equiv complex-cone\ (-a)\ a$

lemma *zero-in-complex-cone* [*simp, intro*]: $a \leq b \implies 0 \in complex-cone\ a\ b$
by (auto *simp: complex-cone-def*)

lemma *complex-coneE*:

```

assumes  $z \in complex-cone\ a\ b$ 
obtains  $r\ \alpha$  where  $r \geq 0\ \alpha \in \{a..b\}\ z = rcis\ r\ \alpha$ 
proof -
  from assms obtain  $y$  where  $y \in \{a..b\}\ z = rcis\ (norm\ z)\ y$ 
  unfolding complex-cone-def by auto
  thus ?thesis using that[of norm z y] by auto
qed

```

lemma *arg-cis* [*simp*]:

```

assumes  $x \in \{-\pi <.. \pi\}$ 
shows  $Arg\ (cis\ x) = x$ 
using assms by (intro cis-Arg-unique) auto

```

lemma *arg-mult-of-real-left* [*simp*]:

```

assumes  $r > 0$ 
shows  $Arg\ (of-real\ r * z) = Arg\ z$ 
proof (cases  $z = 0$ )
  case False
  thus ?thesis
    using Arg-bounded[of z] assms
    by (intro cis-Arg-unique) (auto simp: sgn-mult sgn-of-real cis-Arg)
qed auto

```

lemma *arg-mult-of-real-right* [*simp*]:
assumes $r > 0$
shows $\text{Arg } (z * \text{of-real } r) = \text{Arg } z$
by (*subst mult.commute, subst arg-mult-of-real-left*) (*simp-all add: assms*)

lemma *arg-rcis* [*simp*]:
assumes $x \in \{-\pi < .. \pi\}$ $r > 0$
shows $\text{Arg } (\text{rcis } r x) = x$
using *assms* **by** (*simp add: rcis-def*)

lemma *rcis-in-complex-cone* [*intro*]:
assumes $\alpha \in \{a..b\}$ $r \geq 0$
shows $\text{rcis } r \alpha \in \text{complex-cone } a b$
using *assms* **by** (*auto simp: complex-cone-def*)

lemma *arg-imp-in-complex-cone*:
assumes $\text{Arg } z \in \{a..b\}$
shows $z \in \text{complex-cone } a b$
proof –
have $z = \text{rcis } (\text{norm } z) (\text{Arg } z)$
by (*simp add: rcis-cmod-Arg*)
also have $\dots \in \text{complex-cone } a b$
using *assms* **by** *auto*
finally show *?thesis* .
qed

lemma *complex-cone-altdef*:
assumes $-\pi < a \leq b \leq \pi$
shows $\text{complex-cone } a b = \text{insert } 0 \{z. \text{Arg } z \in \{a..b\}\}$
proof (*intro equalityI subsetI*)
fix z **assume** $z \in \text{complex-cone } a b$
then obtain $r \alpha$ **where** $*$: $r \geq 0 \alpha \in \{a..b\} z = \text{rcis } r \alpha$
by (*auto elim: complex-coneE*)
have $\text{Arg } z \in \{a..b\}$ **if** [*simp*]: $z \neq 0$
proof –
have $r > 0$ **using** *that* $*$ **by** (*subst (asm) **) *auto*
hence $\alpha \in \{a..b\}$
using $*(1,2)$ *assms* **by** (*auto simp: *(1)*)
moreover from *assms* $*(2)$ **have** $\alpha \in \{-\pi < .. \pi\}$
by *auto*
ultimately show *?thesis* **using** $*(3)$ $\langle r > 0 \rangle$
by (*subst **) *auto*
qed
thus $z \in \text{insert } 0 \{z. \text{Arg } z \in \{a..b\}\}$
by *auto*
qed (*use assms in* $\langle \text{auto intro: arg-imp-in-complex-cone} \rangle$)

lemma *nonneg-of-real-in-complex-cone* [*simp, intro*]:

assumes $x \geq 0 \ a \leq 0 \ 0 \leq b$
shows *of-real* $x \in \text{complex-cone } a \ b$
proof –
from *assms* **have** *rcis* $x \ 0 \in \text{complex-cone } a \ b$
by (*intro rcis-in-complex-cone*) *auto*
thus *?thesis* **by** *simp*
qed

lemma *one-in-complex-cone* [*simp, intro*]: $a \leq 0 \implies 0 \leq b \implies 1 \in \text{complex-cone } a \ b$
using *nonneg-of-real-in-complex-cone[of 1]* **by** (*simp del: nonneg-of-real-in-complex-cone*)

lemma *of-nat-in-complex-cone* [*simp, intro*]: $a \leq 0 \implies 0 \leq b \implies \text{of-nat } n \in \text{complex-cone } a \ b$
using *nonneg-of-real-in-complex-cone[of real n]* **by** (*simp del: nonneg-of-real-in-complex-cone*)

2.3 Another integral representation of the Beta function

lemma *complex-cone-inter-nonpos-Reals*:
assumes $-\pi < a \ a \leq b \ b < \pi$
shows $\text{complex-cone } a \ b \cap \mathbb{R}_{\leq 0} = \{0\}$
proof (*safe elim!: nonpos-Reals-cases*)
fix $x :: \text{real}$
assume *complex-of-real* $x \in \text{complex-cone } a \ b \ x \leq 0$
hence $\neg(x < 0)$
using *assms* **by** (*intro notI*) (*auto simp: complex-cone-altdef*)
with $\langle x \leq 0 \rangle$ **show** *complex-of-real* $x = 0$ **by** *auto*
qed (*use assms in auto*)

theorem
assumes $a: a > 0$ **and** $b: b > 0$ ($:: \text{real}$)
shows *has-integral-Beta-real'*:
 $((\lambda u. u \ \text{powr } (b - 1) / (1 + u) \ \text{powr } (a + b)) \ \text{has-integral } \text{Beta } a \ b) \ \{0 <..\}$
and *Beta-conv-nn-integral*:
 $\text{Beta } a \ b = (\int^{+} u. \ \text{ennreal } (\text{indicator } \{0 <..\} \ u * u \ \text{powr } (b - 1) / (1 + u) \ \text{powr } (a + b)) \ \partial \text{lborel})$
proof –
define I **where**
 $I = (\int^{+} u. \ \text{ennreal } (\text{indicator } \{0 <..\} \ u * u \ \text{powr } (b - 1) / (1 + u) \ \text{powr } (a + b)) \ \partial \text{lborel})$
have $\text{Gamma } (a + b) > 0 \ \text{Beta } a \ b > 0$
using *assms* **by** (*simp-all add: add-pos-pos Beta-def*)
from $a \ b$ **have** $\text{ennreal } (\text{Gamma } a * \text{Gamma } b) =$
 $(\int^{+} t. \ \text{ennreal } (\text{indicator } \{0..\} \ t * t \ \text{powr } (a - 1) / \exp t) \ \partial \text{lborel}) *$
 $(\int^{+} t. \ \text{ennreal } (\text{indicator } \{0..\} \ t * t \ \text{powr } (b - 1) / \exp t) \ \partial \text{lborel})$
by (*subst ennreal-mult'*) (*simp-all add: Gamma-conv-nn-integral-real*)
also have $\dots = (\int^{+} t. \ \int^{+} u. \ \text{ennreal } (\text{indicator } \{0..\} \ t * t \ \text{powr } (a - 1) / \exp t) *$
 $\text{ennreal } (\text{indicator } \{0..\} \ u * u \ \text{powr } (b - 1) / \exp u) \ \partial \text{lborel})$

∂lborel
by (*simp add: nn-integral-cmult nn-integral-multc*)
also have $\dots = (\int^{+t}. \text{indicator } \{0<..\} t * (\int^{+u}. \text{indicator } \{0..\} u * t \text{ powr } (a - 1) * u \text{ powr } (b - 1))$
 $\quad / \text{exp } (t + u) \partial\text{lborel}) \partial\text{lborel}$
by (*intro nn-integral-cong-AE AE-I[of - - {0}]*)
(auto simp: indicator-def divide-ennreal ennreal-mult' [symmetric] exp-add mult-ac)
also have $\dots = (\int^{+t}. \text{indicator } \{0<..\} t * (\int^{+u}. \text{indicator } \{0..\} u * t \text{ powr } (a - 1) * u \text{ powr } (b - 1))$
 $\quad / \text{exp } (t + u)$
 $\quad \partial(\text{density } (\text{distr lborel borel } ((* t)) (\lambda x. \text{ennreal } |t|))) \partial\text{lborel}) \partial\text{lborel}$
by (*intro nn-integral-cong mult-indicator-cong, subst lborel-distr-mult' [symmetric]*)
auto
also have $\dots = (\int^{+(t::\text{real})}. \text{indicator } \{0<..\} t * (\int^{+u}. \text{indicator } \{0..\} (u * t) * t \text{ powr } a * (u * t) \text{ powr } (b - 1) / \text{exp } (t + t * u) \partial\text{lborel}) \partial\text{lborel})$
by (*intro nn-integral-cong mult-indicator-cong*)
(auto simp: nn-integral-density nn-integral-distr algebra-simps powr-diff simp flip: ennreal-mult)
also have $\dots = (\int^{+(t::\text{real})}. \int^{+u}. \text{indicator } (\{0<..\} \times \{0..\}) (t, u) * t \text{ powr } a * (u * t) \text{ powr } (b - 1) / \text{exp } (t * (u + 1)) \partial\text{lborel}) \partial\text{lborel}$
by (*subst nn-integral-cmult [symmetric], simp, intro nn-integral-cong*)
(auto simp: indicator-def zero-le-mult-iff algebra-simps)
also have $\dots = (\int^{+(t::\text{real})}. \int^{+u}. \text{indicator } (\{0<..\} \times \{0..\}) (t, u) * t \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{exp } (t * (u + 1)) \partial\text{lborel}) \partial\text{lborel}$
by (*intro nn-integral-cong*) *(auto simp: powr-add powr-diff indicator-def powr-mult field-simps)*
also have $\dots = (\int^{+(u::\text{real})}. \int^{+t}. \text{indicator } (\{0<..\} \times \{0..\}) (t, u) * t \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{exp } (t * (u + 1)) \partial\text{lborel}) \partial\text{lborel}$
by (*rule lborel-pair.Fubini'*) *auto*
also have $\dots = (\int^{+(u::\text{real})}. \text{indicator } \{0..\} u * (\int^{+t}. \text{indicator } \{0<..\} t * t \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{exp } (t * (u + 1)) \partial\text{lborel}) \partial\text{lborel}$
by (*intro nn-integral-cong mult-indicator-cong*) *(auto simp: indicator-def)*
also have $\dots = (\int^{+(u::\text{real})}. \text{indicator } \{0<..\} u * (\int^{+t}. \text{indicator } \{0<..\} t * t \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{exp } (t * (u + 1)) \partial\text{lborel}) \partial\text{lborel}$
by (*intro nn-integral-cong-AE AE-I[of - - {0}]*) *(auto simp: indicator-def)*
also have $\dots = (\int^{+(u::\text{real})}. \text{indicator } \{0<..\} u * (\int^{+t}. \text{indicator } \{0<..\} t * t \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{exp } (t * (u + 1))$
 $\quad \partial(\text{density } (\text{distr lborel borel } ((* (1/(1+u)))) (\lambda x. \text{ennreal } |1/(1+u)|))) \partial\text{lborel}) \partial\text{lborel}$
by (*intro nn-integral-cong mult-indicator-cong, subst lborel-distr-mult' [symmetric]*)
auto
also have $\dots = (\int^{+(u::\text{real})}. \text{indicator } \{0<..\} u * (\int^{+t}. \text{ennreal } (1 / (u + 1)) * \text{ennreal } (\text{indicator } \{0<..\} (t / (u$

$+ 1)) *$
 $(t / (1+u)) \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{ exp } t$
 $\partial \text{lborel } \partial \text{lborel}$
by (*intro nn-integral-cong mult-indicator-cong*)
(auto simp: nn-integral-distr nn-integral-density add-ac)
also have $\dots = (\int^{+u}. \int^{+t}. \text{ indicator } (\{0<..\} \times \{0<..\}) (u, t) *$
 $1/(u+1) * (t / (u+1)) \text{ powr } (a + b - 1) * u \text{ powr } (b - 1) / \text{ exp } t$
 $\partial \text{lborel } \partial \text{lborel})$
by (*subst nn-integral-cmult [symmetric], simp, intro nn-integral-cong*)
(auto simp: indicator-def field-simps divide-ennreal simp flip: ennreal-mult
ennreal-mult')
also have $\dots = (\int^{+u}. \int^{+t}. \text{ ennreal } (\text{ indicator } \{0<..\} u * u \text{ powr } (b - 1) / (1$
 $+ u) \text{ powr } (a + b)) *$
 $\text{ ennreal } (\text{ indicator } \{0<..\} t * t \text{ powr } (a + b - 1) / \text{ exp } t)$
 $\partial \text{lborel } \partial \text{lborel})$
by (*intro nn-integral-cong*)
(auto simp: indicator-def powr-add powr-diff powr-divide powr-minus di-
vide-simps add-ac
simp flip: ennreal-mult)
also have $\dots = I * (\int^{+t}. \text{ indicator } \{0<..\} t * t \text{ powr } (a + b - 1) / \text{ exp } t$
 $\partial \text{lborel})$
by (*simp add: nn-integral-cmult nn-integral-multc I-def*)
also have $(\int^{+t}. \text{ indicator } \{0<..\} t * t \text{ powr } (a + b - 1) / \text{ exp } t \partial \text{lborel}) =$
 $\text{ ennreal } (\text{ Gamma } (a + b))$
using *assms*
by (*subst Gamma-conv-nn-integral-real*)
(auto intro!: nn-integral-cong-AE[OF AE-I[of - - {0}]]
simp: indicator-def split: if-splits)
finally have $\text{ ennreal } (\text{ Gamma } a * \text{ Gamma } b) = I * \text{ ennreal } (\text{ Gamma } (a + b)) .$
hence $\text{ ennreal } (\text{ Gamma } a * \text{ Gamma } b) / \text{ ennreal } (\text{ Gamma } (a + b)) =$
 $I * \text{ ennreal } (\text{ Gamma } (a + b)) / \text{ ennreal } (\text{ Gamma } (a + b))$ **by** *simp*
also have $\dots = I$
using $\langle \text{ Gamma } (a + b) > 0 \rangle$ **by** (*intro ennreal-mult-divide-eq*) (*auto simp:*)
also have $\text{ ennreal } (\text{ Gamma } a * \text{ Gamma } b) / \text{ ennreal } (\text{ Gamma } (a + b)) =$
 $\text{ ennreal } (\text{ Gamma } a * \text{ Gamma } b / \text{ Gamma } (a + b))$
using *assms* **by** (*intro divide-ennreal*) *auto*
also have $\dots = \text{ ennreal } (\text{ Beta } a b)$
by (*simp add: Beta-def*)
finally show $*$: $\text{ ennreal } (\text{ Beta } a b) = I .$

define f **where** $f = (\lambda u. u \text{ powr } (b - 1) / (1 + u) \text{ powr } (a + b))$
have $(\lambda u. \text{ indicator } \{0<..\} u * f u) \text{ has-integral } \text{ Beta } a b$ *UNIV*
using $*$ $\langle \text{ Beta } a b > 0 \rangle$
by (*subst has-integral-iff-nn-integral-lebesgue*)
(auto simp: f-def measurable-completion nn-integral-completion I-def mult-ac)
also have $(\lambda u. \text{ indicator } \{0<..\} u * f u) = (\lambda u. \text{ if } u \in \{0<..\} \text{ then } f u \text{ else } 0)$
by (*auto simp: fun-eq-iff*)
also have $(\dots \text{ has-integral } \text{ Beta } a b) \text{ UNIV} \longleftrightarrow (f \text{ has-integral } \text{ Beta } a b) \{0<..\}$
by (*rule has-integral-restrict-UNIV*)

finally show ... by (simp add: f-def)
qed

lemma has-integral-Beta2:

fixes a :: real

assumes a < -1/2

shows ((λx. (1 + x ^ 2) powr a) has-integral Beta (- a - 1 / 2) (1 / 2) / 2) {0<..}

proof -

define f where f = (λu. u powr (-1/2) / (1 + u) powr (-a))

define C where C = Beta (-a-1/2) (1/2)

have I: (f has-integral C) {0<..}

using has-integral-Beta-real'[of -a-1/2 1/2] assms

by (simp-all add: diff-divide-distrib f-def C-def)

define g where g = (λx. x ^ 2 :: real)

have bij: bij-betw g {0<..} {0<..}

by (intro bij-betwI[of - - - sqrt]) (auto simp: g-def)

have (f absolutely-integrable-on g ' {0<..} ∧ integral (g ' {0<..}) f = C)

using I bij by (simp add: bij-betw-def has-integral-iff absolutely-integrable-on-def f-def)

also have ?this ⟷ ((λx. |2 * x| *_R f (g x)) absolutely-integrable-on {0<..} ∧ integral {0<..} (λx. |2 * x| *_R f (g x)) = C)

using bij by (intro has-absolute-integral-change-of-variables-1' [symmetric]) (auto intro!: derivative-eq-intros simp: g-def bij-betw-def)

finally have ((λx. |2 * x| * f (g x)) has-integral C) {0<..}

by (simp add: absolutely-integrable-on-def f-def has-integral-iff)

also have ?this ⟷ ((λx::real. 2 * (1 + x²) powr a) has-integral C) {0<..}

by (intro has-integral-cong) (auto simp: f-def g-def powr-def exp-minus ln-realpow field-simps)

finally have ((λx::real. 1/2 * (2 * (1 + x²) powr a)) has-integral 1/2 * C) {0<..}

by (intro has-integral-mult-right)

thus ?thesis by (simp add: C-def)

qed

lemma has-integral-Beta3:

fixes a b :: real

assumes a < -1/2 and b > 0

shows ((λx. (b + x ^ 2) powr a) has-integral

Beta (-a - 1/2) (1/2) / 2 * b powr (a + 1/2)) {0<..}

proof -

define C where C = Beta (- a - 1 / 2) (1 / 2) / 2

have int: nn-integral lborel (λx. indicator {0<..} x * (1 + x ^ 2) powr a) = C

using nn-integral-has-integral-lebesgue[OF - has-integral-Beta2[OF assms(1)]]

by (auto simp: C-def)

have nn-integral lborel (λx. indicator {0<..} x * (b + x ^ 2) powr a) =

(∫⁺x. ennreal (indicat-real {0<..} (x * sqrt b) * (b + (x * sqrt b)²) powr a

```

* sqrt b) ∂lborel)
  using assms
  by (subst lborel-distr-mult'[of sqrt b])
    (auto simp: nn-integral-density nn-integral-distr mult-ac simp flip: ennreal-mult)
  also have ... = (∫+x. ennreal (indicat-real {0<..} x * (b * (1 + x2))) powr
a * sqrt b) ∂lborel)
  using assms
  by (intro nn-integral-cong) (auto simp: indicator-def field-simps zero-less-mult-iff)
  also have ... = (∫+x. ennreal (indicat-real {0<..} x * b powr (a + 1/2) * (1
+ x2) powr a) ∂lborel)
  using assms
  by (intro nn-integral-cong) (auto simp: indicator-def powr-add powr-half-sqrt
powr-mult)
  also have ... = b powr (a + 1/2) * (∫+x. ennreal (indicat-real {0<..} x * (1
+ x2) powr a) ∂lborel)
  using assms by (subst nn-integral-cmult [symmetric]) (simp-all add: mult-ac
flip: ennreal-mult)
  also have (∫+x. ennreal (indicat-real {0<..} x * (1 + x2) powr a) ∂lborel)
= C
  using int by simp
  also have ennreal (b powr (a + 1/2)) * ennreal C = ennreal (C * b powr (a +
1/2))
  using assms by (subst ennreal-mult) (auto simp: C-def mult-ac Beta-def)
  finally have *: (∫+x. ennreal (indicat-real {0<..} x * (b + x2) powr a) ∂lborel)
= ... .
  hence ((λx. indicat {0<..} x * (b + x2) powr a) has-integral C * b powr (a
+ 1/2)) UNIV
  using assms
  by (subst has-integral-iff-nn-integral-lebesgue)
    (auto simp: C-def measurable-completion nn-integral-completion Beta-def)
  also have (λx. indicat {0<..} x * (b + x2) powr a) =
    (λx. if x ∈ {0<..} then (b + x2) powr a else 0)
  by (auto simp: fun-eq-iff)
  finally show ?thesis
  by (subst (asm) has-integral-restrict-UNIV) (auto simp: C-def)
qed

```

2.4 Asymptotics of the real $\log \Gamma$ function and its derivatives

This is the error term that occurs in the expansion of $\ln\text{-Gamma}$. It can be shown to be of order $O(s^{-n})$.

definition *stirling-integral* :: $\text{nat} \Rightarrow 'a :: \{\text{real-normed-div-algebra, banach}\} \Rightarrow 'a$
where

$$\text{stirling-integral } n \ s = \lim (\lambda N. \text{integral } \{0..N\} (\lambda x. \text{of-real } (\text{pbernpoly } n \ x) / (\text{of-real } (x + s)^n)))$$

context

fixes $s :: \text{complex}$ **assumes** $s: s \notin \mathbb{R}_{\leq 0}$

fixes *approx* :: $\text{nat} \Rightarrow \text{complex}$

defines $approx \equiv (\lambda N.$
 $(\sum n = 1..<N. s / of-nat n - \ln (1 + s / of-nat n)) - (euler-mascheroni * s$
 $+ \ln s) - \longrightarrow \ln\text{-Gamma } s$
 $(\ln\text{-Gamma } (of-nat N) - \ln (2 * pi / of-nat N) / 2 - of-nat N * \ln (of-nat$
 $N) + of-nat N) - \longrightarrow 0$
 $s * (harm (N - 1) - \ln (of-nat (N - 1)) - euler-mascheroni) + \longrightarrow 0$
 $s * (\ln (of-nat N + s) - \ln (of-nat (N - 1))) - \longrightarrow 0$
 $(1/2) * (\ln (of-nat N + s) - \ln (of-nat N)) + \longrightarrow 0$
 $of-nat N * (\ln (of-nat N + s) - \ln (of-nat N)) - \longrightarrow s$
 $(s - 1/2) * \ln s - \ln (2 * pi) / 2)$

begin

qualified lemma

assumes $N: N > 0$

shows *integrable-pbernpoly-1:*

$(\lambda x. of-real (-pbernpoly 1 x) / (of-real x + s))$ *integrable-on* $\{0..real N\}$

and *integral-pbernpoly-1-aux:*

integral $\{0..real N\}$ $(\lambda x. -of-real (pbernpoly 1 x) / (of-real x + s)) =$

$approx N$

and *has-integral-pbernpoly-1:*

$((\lambda x. pbernpoly 1 x / (x + s))$ *has-integral*
 $(\sum m < N. (of-nat m + 1 / 2 + s) * (\ln (of-nat m + s) -$
 $\ln (of-nat m + 1 + s)) + 1))$ $\{0..real N\}$

proof –

let $?A = (\lambda n. \{of-nat n..of-nat (n+1)\}) ' \{0..<N\}$

have *has-integral:*

$((\lambda x. -pbernpoly 1 x / (x + s))$ *has-integral*
 $(of-nat n + 1/2 + s) * (\ln (of-nat (n + 1) + s) - \ln (of-nat n + s))$

– 1)

$\{of-nat n..of-nat (n + 1)\}$ **for** n

proof (*rule has-integral-spike*)

have $((\lambda x. (of-nat n + 1/2 + s) * (1 / (of-real x + s)) - 1)$ *has-integral*
 $(of-nat n + 1/2 + s) * (\ln (of-real (real (n + 1)) + s) - \ln (of-real$
 $(real n) + s)) - 1)$

$\{of-nat n..of-nat (n + 1)\}$

using s *has-integral-const-real*[*of 1 of-nat n of-nat (n + 1)*]

by (*intro has-integral-diff has-integral-mult-right fundamental-theorem-of-calculus*)

(*auto intro!:* *derivative-eq-intros has-vector-derivative-real-field*

simp: *has-field-derivative-iff-has-vector-derivative [symmetric] field-simps*
complex-nonpos-Reals-iff)

thus $((\lambda x. (of-nat n + 1/2 + s) * (1 / (of-real x + s)) - 1)$ *has-integral*
 $(of-nat n + 1/2 + s) * (\ln (of-nat (n + 1) + s) - \ln (of-nat n + s))$

– 1)

$\{of-nat n..of-nat (n + 1)\}$ **by** *simp*

show $-pbernpoly 1 x / (x + s) = (of-nat n + 1/2 + s) * (1 / (x + s)) - 1$

if $x \in \{of-nat n..of-nat (n + 1)\} - \{of-nat (n + 1)\}$ **for** x

proof –

have $x: x \geq real n$ $x < real (n + 1)$ **using** *that by simp-all*

hence $\text{floor } x = \text{int } n$ **by** *linarith*
moreover from $s \ x$ **have** *complex-of-real* $x \neq -s$
by (*auto simp add: complex-eq-iff complex-nonpos-Reals-iff simp del: of-nat-Suc*)
ultimately show $-p\text{bernpoly } 1 \ x / (x + s) = (\text{of-nat } n + 1/2 + s) * (1 / (x + s)) - 1$
by (*auto simp: pbernpoly-def bernpoly-def frac-def divide-simps add-eq-0-iff2*)
qed
qed *simp-all*
hence $*$: $\bigwedge I. I \in ?A \implies ((\lambda x. -p\text{bernpoly } 1 \ x / (x + s)) \text{ has-integral } (\text{Inf } I + 1/2 + s) * (\ln (\text{Inf } I + 1 + s) - \ln (\text{Inf } I + s)) - 1) \ I$
by (*auto simp: add-ac*)
have $((\lambda x. -p\text{bernpoly } 1 \ x / (x + s)) \text{ has-integral } (\sum I \in ?A. (\text{Inf } I + 1 / 2 + s) * (\ln (\text{Inf } I + 1 + s) - \ln (\text{Inf } I + s)) - 1))$
 $(\bigcup n \in \{0..<N\}. \{\text{real } n.. \text{real } (n + 1)\})$ **(is** $(- \text{ has-integral } ?i) -$
apply (*intro has-integral-Union * finite-imageI*)
apply (*force intro!: negligible-atLeastAtMostI pairwiseI*)
done
hence *has-integral*: $((\lambda x. -p\text{bernpoly } 1 \ x / (x + s)) \text{ has-integral } ?i) \{0.. \text{real } N\}$
by (*subst has-integral-spike-set-eq*)
(use Union-atLeastAtMost assms in <auto simp: intro!: empty-imp-negligible>)
hence $(\lambda x. -p\text{bernpoly } 1 \ x / (x + s))$ *integrable-on* $\{0.. \text{real } N\}$
and *integral*: $\text{integral } \{0.. \text{real } N\} (\lambda x. -p\text{bernpoly } 1 \ x / (x + s)) = ?i$
by (*simp-all add: has-integral-iff*)
show $(\lambda x. -p\text{bernpoly } 1 \ x / (x + s))$ *integrable-on* $\{0.. \text{real } N\}$ **by fact**

note *has-integral-neg* [*OF has-integral*]
also have $-?i = (\sum x < N. (\text{of-nat } x + 1 / 2 + s) * (\ln (\text{of-nat } x + s) - \ln (\text{of-nat } x + 1 + s))) + 1$
by (*subst sum.reindex*)
(simp-all add: inj-on-def atLeast0LessThan algebra-simps sum-negf [symmetric])
finally show *has-integral*:
 $((\lambda x. \text{of-real } (p\text{bernpoly } 1 \ x) / (\text{of-real } x + s)) \text{ has-integral } \dots) \{0.. \text{real } N\}$ **by**
simp

note *integral*
also have $?i = (\sum n < N. (\text{of-nat } n + 1 / 2 + s) * (\ln (\text{of-nat } n + 1 + s) - \ln (\text{of-nat } n + s))) - N$ **(is** $- = ?S - -$)
by (*subst sum.reindex*) (*simp-all add: inj-on-def sum-subtractf atLeast0LessThan*)
also have $?S = (\sum n < N. \text{of-nat } n * (\ln (\text{of-nat } n + 1 + s) - \ln (\text{of-nat } n + s))) +$
 $(s + 1 / 2) * (\sum n < N. \ln (\text{of-nat } (\text{Suc } n) + s) - \ln (\text{of-nat } n + s))$
(is $- = ?S1 + - * ?S2$) **by** (*simp add: algebra-simps sum.distrib sum-subtractf sum-distrib-left*)
also have $?S2 = \ln (\text{of-nat } N + s) - \ln s$ **by** (*subst sum-lessThan-telescope*)
simp
also have $?S1 = (\sum n = 1..<N. \text{of-nat } n * (\ln (\text{of-nat } n + 1 + s) - \ln (\text{of-nat } n + s)))$

by (*intro sum.mono-neutral-right*) *auto*
 also have ... = $(\sum_{n=1..<N}. \text{of-nat } n * \ln (\text{of-nat } n + 1 + s)) - (\sum_{n=1..<N}. \text{of-nat } n * \ln (\text{of-nat } n + s))$
 by (*simp add: algebra-simps sum-subtractf*)
 also have $(\sum_{n=1..<N}. \text{of-nat } n * \ln (\text{of-nat } n + 1 + s)) =$
 $(\sum_{n=1..<N}. (\text{of-nat } n - 1) * \ln (\text{of-nat } n + s)) + (N - 1) * \ln (\text{of-nat } N + s)$
 by (*induction N*) (*simp-all add: add-ac of-nat-diff*)
 also have ... - $(\sum_{n=1..<N}. \text{of-nat } n * \ln (\text{of-nat } n + s)) =$
 $-(\sum_{n=1..<N}. \ln (\text{of-nat } n + s)) + (N - 1) * \ln (\text{of-nat } N + s)$
 by (*induction N*) (*simp-all add: algebra-simps*)
 also from *s* have *neg: s + of-nat x ≠ 0 for x*
 by (*auto simp: complex-nonpos-Reals-iff complex-eq-iff*)
 hence $(\sum_{n=1..<N}. \ln (\text{of-nat } n + s)) = (\sum_{n=1..<N}. \ln (\text{of-nat } n)) + \ln (1 + s/n)$
 by (*intro sum.cong refl, subst Ln-times-of-nat [symmetric]*) (*auto simp: divide-simps add-ac*)
 also have ... = $\ln (\text{fact } (N - 1)) + (\sum_{n=1..<N}. \ln (1 + s/n))$
 by (*induction N*) (*simp-all add: Ln-times-of-nat fact-reduce add-ac*)
 also have $(\sum_{n=1..<N}. \ln (1 + s/n)) = -(\sum_{n=1..<N}. s / n - \ln (1 + s/n))$
 $+ s * (\sum_{n=1..<N}. 1 / \text{of-nat } n)$
 by (*simp add: sum-distrib-left sum-subtractf*)
 also from *N* have $\ln (\text{fact } (N - 1)) = \ln\text{-Gamma } (\text{of-nat } N :: \text{complex})$
 by (*simp add: ln-Gamma-complex-conv-fact*)
 also have $\{1..<N\} = \{1..N - 1\}$ by *auto*
 hence $(\sum_{n=1..<N}. 1 / \text{of-nat } n) = (\text{harm } (N - 1) :: \text{complex})$
 by (*simp add: harm-def divide-simps*)
 also have $-(\ln\text{-Gamma } (\text{of-nat } N)) + (- (\sum_{n=1..<N}. s / \text{of-nat } n - \ln (1 + s / \text{of-nat } n)) +$
 $s * \text{harm } (N - 1))) + \text{of-nat } (N - 1) * \ln (\text{of-nat } N + s) +$
 $(s + 1 / 2) * (\ln (\text{of-nat } N + s) - \ln s) - \text{of-nat } N = \text{approx } N$
 using *N* by (*simp add: field-simps of-nat-diff ln-div approx-def Ln-of-nat ln-Gamma-complex-of-real [symmetric]*)
 finally show *integral* $\{0.. \text{of-nat } N\} (\lambda x. -\text{of-real } (\text{pbernpoly } 1 x) / (\text{of-real } x + s)) = \dots$
 by *simp*
 qed

lemma *integrable-ln-Gamma-aux:*

shows $(\lambda x. \text{of-real } (\text{pbernpoly } n x) / (\text{of-real } x + s) ^ n)$ *integrable-on* $\{0.. \text{real } N\}$

proof (*cases n = 1*)

case *True*

with *s* **show** *?thesis* using *integrable-neg[OF integrable-pbernpoly-1[of N]]*

by (*cases N = 0*) (*simp-all add: integrable-negligible*)

next

case *False*

from *s* **have** *of-real x + s ≠ 0 if x ≥ 0 for x* using *that*

by (*auto simp: complex-eq-iff add-eq-0-iff2 complex-nonpos-Reals-iff*)

with *False s show ?thesis*
by (*auto intro! integrable-continuous-real continuous-intros*)
qed

This following proof is based on “Rudiments of the theory of the gamma function” by Bruce Berndt [1].

lemma *tendsto-of-real-0-I*:
 $(f \longrightarrow 0) G \implies ((\lambda x. (\text{of-real } (f x))) \longrightarrow (0 :: 'a :: \text{real-normed-div-algebra}))$
 G
using *tendsto-of-real-iff by force*

qualified lemma *integral-pbernpoly-1*:
 $(\lambda N. \text{integral } \{0..real N\} (\lambda x. \text{pbernpoly } 1 x / (x + s)))$
 $\longrightarrow -\ln\text{-Gamma } s - s + (s - 1 / 2) * \ln s + \ln (2 * \text{pi}) / 2$

proof –

have *neg: s + of-real x ≠ 0 if x ≥ 0 for x :: real*
using *that s by (auto simp: complex-eq-iff complex-nonpos-Reals-iff)*
have (*approx* $\longrightarrow \ln\text{-Gamma } s - 0 - 0 + 0 - 0 + s - (s - 1/2) * \ln s - \ln (2 * \text{pi}) / 2$) *at-top*
unfolding *approx-def*
proof (*intro tendsto-add tendsto-diff*)
from *s have s': s ∉ ℤ_{≤0} by (auto simp: complex-nonpos-Reals-iff elim!: non-pos-Ints-cases)*
have ($\lambda n. \sum_{i=1..<n.} s / \text{of-nat } i - \ln (1 + s / \text{of-nat } i)$) \longrightarrow
 $\ln\text{-Gamma } s + \text{euler-mascheroni} * s + \ln s$ (**is** *?f* \longrightarrow -)
using *ln-Gamma-series'-aux[OF s'] unfolding sums-def*
by (*subst filterlim-sequentially-Suc [symmetric], subst (asm) sum.atLeast1-atMost-eq [symmetric]*)
(simp add: atLeastLessThanSuc-atLeastAtMost)
thus ($\lambda n. ?f n - (\text{euler-mascheroni} * s + \ln s)$) $\longrightarrow \ln\text{-Gamma } s$) *at-top*
by (*auto intro: tendsto-eq-intros*)
next
show ($\lambda x. \text{complex-of-real } (\ln\text{-Gamma } (\text{real } x) - \ln (2 * \text{pi} / \text{real } x) / 2 - \text{real } x * \ln (\text{real } x) + \text{real } x)$) $\longrightarrow 0$
proof (*intro tendsto-of-real-0-I*
filterlim-compose[OF tendsto-sandwich filterlim-real-sequentially])
show *eventually* ($\lambda x :: \text{real. } \ln\text{-Gamma } x - \ln (2 * \text{pi} / x) / 2 - x * \ln x + x$
 ≥ 0) *at-top*
using *eventually-ge-at-top[of 1::real]*
by *eventually-elim (insert ln-Gamma-bounds(1), simp add: algebra-simps)*
show *eventually* ($\lambda x :: \text{real. } \ln\text{-Gamma } x - \ln (2 * \text{pi} / x) / 2 - x * \ln x + x$
 \leq
 $1 / 12 * \text{inverse } x$) *at-top*
using *eventually-ge-at-top[of 1::real]*
by *eventually-elim (insert ln-Gamma-bounds(2), simp add: field-simps)*
show ($\lambda x :: \text{real. } 1 / 12 * \text{inverse } x \longrightarrow 0$) *at-top*
by (*intro tendsto-mult-right-zero tendsto-inverse-0-at-top filterlim-ident*)
qed *simp-all*
next

```

have ( $\lambda x. s * \text{of-real} (\text{harm} (x - 1) - \ln (\text{real} (x - 1)) - \text{euler-mascheroni})$ )
 $\longrightarrow$ 
   $s * \text{of-real} (\text{euler-mascheroni} - \text{euler-mascheroni})$ 
by (subst filterlim-sequentially-Suc [symmetric], intro tendsto-intros)
  (insert euler-mascheroni-LIMSEQ, simp-all)
also have  $?this \longleftrightarrow (\lambda x. s * (\text{harm} (x - 1) - \ln (\text{of-nat} (x - 1)) - \text{euler-mascheroni})) \longrightarrow 0$ 
by (intro filterlim-cong refl eventually-mono[OF eventually-gt-at-top[of 1::nat]])

  (auto simp: Ln-of-nat of-real-harm)
finally show ( $\lambda x. s * (\text{harm} (x - 1) - \ln (\text{of-nat} (x - 1)) - \text{euler-mascheroni})$ )
 $\longrightarrow 0$  .
next
have ( $(\lambda x. \ln (1 + (s + 1) / \text{of-real} x)) \longrightarrow \ln (1 + 0)$ ) at-top (is  $?P$ )
by (intro tendsto-intros tendsto-divide-0[OF tendsto-const])
  (simp-all add: filterlim-ident filterlim-at-infinity-conv-norm-at-top filterlim-abs-real)
also have  $\ln (\text{of-real} (x + 1) + s) - \ln (\text{complex-of-real} x) = \ln (1 + (s + 1) / \text{of-real} x)$ 
if  $x > 1$  for  $x$  using that  $s$ 
using Ln-divide-of-real[of x of-real (x + 1) + s, symmetric] neq[of x + 1]
by (simp add: field-simps Ln-of-real)
hence  $?P \longleftrightarrow ((\lambda x. \ln (\text{of-real} (x + 1) + s) - \ln (\text{of-real} x)) \longrightarrow 0)$  at-top
by (intro filterlim-cong refl)
  (auto intro: eventually-mono[OF eventually-gt-at-top[of 1::real]])
finally have ( $(\lambda n. \ln (\text{of-real} (\text{real} n + 1) + s) - \ln (\text{of-real} (\text{real} n))) \longrightarrow 0$ ) at-top
by (rule filterlim-compose[OF - filterlim-real-sequentially])
hence ( $(\lambda n. \ln (\text{of-nat} n + s) - \ln (\text{of-nat} (n - 1))) \longrightarrow 0$ ) at-top
by (subst filterlim-sequentially-Suc [symmetric] (simp add: add-ac))
thus ( $\lambda x. s * (\ln (\text{of-nat} x + s) - \ln (\text{of-nat} (x - 1))) \longrightarrow 0$ )
by (rule tendsto-mult-right-zero)
next
have ( $(\lambda x. \ln (1 + s / \text{of-real} x)) \longrightarrow \ln (1 + 0)$ ) at-top (is  $?P$ )
by (intro tendsto-intros tendsto-divide-0[OF tendsto-const])
  (simp-all add: filterlim-ident filterlim-at-infinity-conv-norm-at-top filterlim-abs-real)
also have  $\ln (\text{of-real} x + s) - \ln (\text{of-real} x) = \ln (1 + s / \text{of-real} x)$  if  $x > 0$ 
for  $x$ 
using Ln-divide-of-real[of x of-real x + s] neq[of x] that
by (auto simp: field-simps Ln-of-real)
hence  $?P \longleftrightarrow ((\lambda x. \ln (\text{of-real} x + s) - \ln (\text{of-real} x)) \longrightarrow 0)$  at-top
using  $s$  by (intro filterlim-cong refl)
  (auto intro: eventually-mono [OF eventually-gt-at-top[of 1::real]])
finally have ( $\lambda x. (1/2) * (\ln (\text{of-real} (\text{real} x) + s) - \ln (\text{of-real} (\text{real} x)))$ )
 $\longrightarrow 0$ 
by (rule tendsto-mult-right-zero[OF filterlim-compose[OF - filterlim-real-sequentially]])
thus ( $\lambda x. (1/2) * (\ln (\text{of-nat} x + s) - \ln (\text{of-nat} x)) \longrightarrow 0$ ) by simp
next

```

have $((\lambda x. x * (\ln (1 + s / \text{of-real } x))) \longrightarrow s)$ *at-top* (**is** ?P)
by (*rule stirling-limit-aux2*)
also have $\ln (1 + s / \text{of-real } x) = \ln (\text{of-real } x + s) - \ln (\text{of-real } x)$ **if** $x > 1$
for x
using *that* s *Ln-divide-of-real* [*of* x *of-real* $x + s$, *symmetric*] *neq*[*of* x]
by (*auto simp: Ln-of-real field-simps*)
hence ?P $\longleftrightarrow ((\lambda x. \text{of-real } x * (\ln (\text{of-real } x + s) - \ln (\text{of-real } x))) \longrightarrow s)$
at-top
by (*intro filterlim-cong refl*)
(auto intro: eventually-mono[OF eventually-gt-at-top[of 1::real]])
finally have $(\lambda n. \text{of-real } (\text{real } n) * (\ln (\text{of-real } (\text{real } n) + s) - \ln (\text{of-real } (\text{real } n)))) \longrightarrow s$
by (*rule filterlim-compose[OF - filterlim-real-sequentially]*)
thus $(\lambda n. \text{of-nat } n * (\ln (\text{of-nat } n + s) - \ln (\text{of-nat } n))) \longrightarrow s$ **by** *simp*
qed *simp-all*
also have ?this $\longleftrightarrow ((\lambda N. \text{integral } \{0.. \text{real } N\} (\lambda x. -\text{pbernpoly } 1 \ x / (x + s))) \longrightarrow$
 \longrightarrow
 $\ln\text{-Gamma } s + s - (s - 1/2) * \ln s - \ln (2 * \text{pi}) / 2)$ *at-top*
using *integral-pbernpoly-1-aux*
by (*intro filterlim-cong refl*)
(auto intro: eventually-mono[OF eventually-gt-at-top[of 0::nat]])
also have $(\lambda N. \text{integral } \{0.. \text{real } N\} (\lambda x. -\text{pbernpoly } 1 \ x / (x + s))) =$
 $(\lambda N. -\text{integral } \{0.. \text{real } N\} (\lambda x. \text{pbernpoly } 1 \ x / (x + s)))$
by (*simp add: fun-eq-iff*)
finally show ?thesis **by** (*simp add: tendsto-minus-cancel-left [symmetric] algebra-simps*)
qed

qualified lemma *pbernpoly-integral-conv-pbernpoly-integral-Suc*:

assumes $n \geq 1$
shows $\text{integral } \{0.. \text{real } N\} (\lambda x. \text{pbernpoly } n \ x / (x + s) ^ n) =$
 $\text{of-real } (\text{pbernpoly } (\text{Suc } n) (\text{real } N)) / (\text{of-nat } (\text{Suc } n) * (s + \text{of-nat } N)$
 $^ n) -$
 $\text{of-real } (\text{bernoulli } (\text{Suc } n)) / (\text{of-nat } (\text{Suc } n) * s ^ n) + \text{of-nat } n / \text{of-nat}$
 $(\text{Suc } n) *$
 $\text{integral } \{0.. \text{real } N\} (\lambda x. \text{of-real } (\text{pbernpoly } (\text{Suc } n) \ x) / (\text{of-real } x +$
 $s) ^ \text{Suc } n)$
proof –
note [*derivative-intros*] = *has-field-derivative-pbernpoly-Suc'*
define I **where** $I = -\text{of-real } (\text{pbernpoly } (\text{Suc } n) (\text{of-nat } N)) / (\text{of-nat } (\text{Suc } n)$
 $* (\text{of-nat } N + s) ^ n) +$
 $\text{of-real } (\text{bernoulli } (\text{Suc } n) / \text{real } (\text{Suc } n)) / s ^ n +$
 $\text{integral } \{0.. \text{real } N\} (\lambda x. \text{of-real } (\text{pbernpoly } n \ x) / (\text{of-real } x + s) ^ n)$
have $((\lambda x. (-\text{of-nat } n * \text{inverse } (\text{of-real } x + s) ^ \text{Suc } n) * (\text{of-real } (\text{pbernpoly } (\text{Suc } n) \ x) / (\text{of-nat } (\text{Suc } n))))$
 $\text{has-integral } -I) \{0.. \text{real } N\}$
proof (*rule integration-by-parts-interior-strong[OF bounded-bilinear-mult]*)
fix $x :: \text{real}$ **assume** $x \in \{0 <.. < \text{real } N\} - \text{real } ' \{0.. N\}$

have $x \notin \mathbb{Z}$
proof
 assume $x \in \mathbb{Z}$
 then obtain n **where** $x = \text{of-int } n$ **by** (*auto elim!: Ints-cases*)
 with x **have** x' : $x = \text{of-nat } (\text{nat } n)$ **by** *simp*
 from x **show** *False* **by** (*auto simp: x'*)
qed
hence $((\lambda x. \text{of-real } (\text{pbernpoly } (\text{Suc } n) x / \text{of-nat } (\text{Suc } n))) \text{ has-vector-derivative } \text{complex-of-real } (\text{pbernpoly } n x)) \text{ (at } x)$
 by (*intro has-vector-derivative-of-real*) (*auto intro!: derivative-eq-intros*)
thus $((\lambda x. \text{of-real } (\text{pbernpoly } (\text{Suc } n) x) / \text{of-nat } (\text{Suc } n)) \text{ has-vector-derivative } \text{complex-of-real } (\text{pbernpoly } n x)) \text{ (at } x)$ **by** *simp*
from x s **have** $\text{complex-of-real } x + s \neq 0$
 by (*auto simp: complex-eq-iff complex-nonpos-Reals-iff*)
thus $((\lambda x. \text{inverse } (\text{of-real } x + s) ^ n) \text{ has-vector-derivative } - \text{of-nat } n * \text{inverse } (\text{of-real } x + s) ^ (\text{Suc } n)) \text{ (at } x)$ **using** x s *assms*
 by (*auto intro!: derivative-eq-intros has-vector-derivative-real-field simp: divide-simps power-add [symmetric]*)
 simp del: power-Suc
next
 have $\text{complex-of-real } x + s \neq 0$ **if** $x \geq 0$ **for** x
 using *that* s **by** (*auto simp: complex-eq-iff complex-nonpos-Reals-iff*)
 thus $\text{continuous-on } \{0.. \text{real } N\} (\lambda x. \text{inverse } (\text{of-real } x + s) ^ n)$
 $\text{continuous-on } \{0.. \text{real } N\} (\lambda x. \text{complex-of-real } (\text{pbernpoly } (\text{Suc } n) x) / \text{of-nat } (\text{Suc } n))$
 using *assms* s **by** (*auto intro!: continuous-intros simp del: of-nat-Suc*)
next
 have $((\lambda x. \text{inverse } (\text{of-real } x + s) ^ n * \text{of-real } (\text{pbernpoly } n x)) \text{ has-integral } \text{pbernpoly } (\text{Suc } n) (\text{of-nat } N) / (\text{of-nat } (\text{Suc } n) * (\text{of-nat } N + s) ^ n) - \text{of-real } (\text{bernoulli } (\text{Suc } n) / \text{real } (\text{Suc } n)) / s ^ n - -I) \{0.. \text{real } N\}$
 using *integrable-ln-Gamma-aux*[*of n N*] *assms*
 by (*auto simp: I-def has-integral-integral divide-simps*)
 thus $((\lambda x. \text{inverse } (\text{of-real } x + s) ^ n * \text{of-real } (\text{pbernpoly } n x)) \text{ has-integral } \text{inverse } (\text{of-real } (\text{real } N) + s) ^ n * (\text{of-real } (\text{pbernpoly } (\text{Suc } n) (\text{real } N)) /$
 $\text{of-nat } (\text{Suc } n)) -$
 $\text{inverse } (\text{of-real } 0 + s) ^ n * (\text{of-real } (\text{pbernpoly } (\text{Suc } n) 0) / \text{of-nat } (\text{Suc } n)) - -I)$
 $\{0.. \text{real } N\}$ **by** (*simp-all add: field-simps*)
qed *simp-all*
 also have $(\lambda x. - \text{of-nat } n * \text{inverse } (\text{of-real } x + s) ^ (\text{Suc } n) * (\text{of-real } (\text{pbernpoly } (\text{Suc } n) x) / \text{of-nat } (\text{Suc } n))) =$
 $(\lambda x. - (\text{of-nat } n / \text{of-nat } (\text{Suc } n)) * \text{of-real } (\text{pbernpoly } (\text{Suc } n) x) / (\text{of-real } x + s) ^ (\text{Suc } n))$
 by (*simp add: divide-simps fun-eq-iff*)
finally have $((\lambda x. - (\text{of-nat } n / \text{of-nat } (\text{Suc } n)) * \text{of-real } (\text{pbernpoly } (\text{Suc } n) x) / (\text{of-real } x + s) ^ (\text{Suc } n)) \text{ has-integral } -I) \{0.. \text{real } N\}$.
from *has-integral-neg*[*OF this*] **show** *?thesis*

by (auto simp add: I-def has-integral-iff algebra-simps integral-mult-right [symmetric])

simp del: power-Suc of-nat-Suc)

qed

lemma *pbernpoly-over-power-tendsto-0*:

assumes $n > 0$

shows $(\lambda x. \text{of-real } (\text{pbernpoly } (\text{Suc } n) (\text{real } x)) / (\text{of-nat } (\text{Suc } n) * (s + \text{of-nat } x) ^ n)) \longrightarrow 0$

proof –

from s have *neg*: $s + \text{of-nat } n \neq 0$ for n

by (auto simp: complex-eq-iff complex-nonpos-Reals-iff)

obtain c where $c: \bigwedge x. \text{norm } (\text{pbernpoly } (\text{Suc } n) x) \leq c$

using *bounded-pbernpoly* by auto

have *eventually* $(\lambda x. \text{real } x + \text{Re } s > 0)$ at-top

by *real-asymp*

hence *eventually* $(\lambda x. \text{norm } (\text{of-real } (\text{pbernpoly } (\text{Suc } n) (\text{real } x)) / (\text{of-nat } (\text{Suc } n) * (s + \text{of-nat } x) ^ n)) \leq (c / \text{real } (\text{Suc } n)) / (\text{real } x + \text{Re } s) ^ n)$ at-top

using *eventually-gt-at-top*[of $0::\text{nat}$]

proof *eventually-elim*

case (*elim* x)

have $\text{norm } (\text{of-real } (\text{pbernpoly } (\text{Suc } n) (\text{real } x)) / (\text{of-nat } (\text{Suc } n) * (s + \text{of-nat } x) ^ n)) \leq (c / \text{real } (\text{Suc } n)) / \text{norm } (s + \text{of-nat } x) ^ n$ (*is* \leq *?rhs*) using c [of x]

by (auto simp: norm-divide norm-mult norm-power neg field-simps simp del: of-nat-Suc)

also have $(\text{real } x + \text{Re } s) \leq \text{cmod } (s + \text{of-nat } x)$

using *complex-Re-le-cmod*[of $s + \text{of-nat } x$] s by (auto simp add: complex-nonpos-Reals-iff)

hence *?rhs* $\leq (c / \text{real } (\text{Suc } n)) / (\text{real } x + \text{Re } s) ^ n$ using s *elim* c [of 0] *neg*[of x]

by (*intro divide-left-mono power-mono mult-pos-pos divide-nonneg-pos zero-less-power*) *auto*

finally show *?case* .

qed

moreover have $(\lambda x. (c / \text{real } (\text{Suc } n)) / (\text{real } x + \text{Re } s) ^ n) \longrightarrow 0$

using $\langle n > 0 \rangle$ by *real-asymp*

ultimately show *?thesis* by (*rule Lim-null-comparison*)

qed

lemma *convergent-stirling-integral*:

assumes $n > 0$

shows *convergent* $(\lambda N. \text{integral } \{0.. \text{real } N\}$

$(\lambda x. \text{of-real } (\text{pbernpoly } n x) / (\text{of-real } x + s) ^ n))$ (*is convergent* (*?f* n))

proof –

have *convergent* (*?f* ($\text{Suc } n$)) for n

proof (*induction* n)

case 0


```

thus ?case using integral-pbernpoly-1 by (auto intro!: convergentI)
next
case (Suc n)
have convergent (λN. ?f (Suc n) N -
  of-real (pbernpoly (Suc (Suc n)) (real N)) /
  (of-nat (Suc (Suc n)) * (s + of-nat N) ^ Suc n) +
  of-real (bernoulli (Suc (Suc n)) / (real (Suc (Suc n)))) / s ^ Suc n)
(is convergent ?g)
by (intro convergent-add convergent-diff Suc
  convergent-const convergentI[OF pbernpoly-over-power-tendsto-0]) simp-all
also have ?g = (λN. of-nat (Suc n) / of-nat (Suc (Suc n)) * ?f (Suc (Suc n))
N) using s
by (subst pbernpoly-integral-conv-pbernpoly-integral-Suc)
  (auto simp: fun-eq-iff field-simps simp del: of-nat-Suc power-Suc)
also have convergent ... ↔ convergent (?f (Suc (Suc n)))
by (intro convergent-mult-const-iff) (simp-all del: of-nat-Suc)
finally show ?case .
qed
from this[of n - 1] assms show ?thesis by simp
qed

lemma stirling-integral-conv-stirling-integral-Suc:
assumes n > 0
shows stirling-integral n s =
  of-nat n / of-nat (Suc n) * stirling-integral (Suc n) s -
  of-real (bernoulli (Suc n)) / (of-nat (Suc n) * s ^ n)

proof -
have (λN. of-real (pbernpoly (Suc n) (real N)) / (of-nat (Suc n) * (s + of-nat
N) ^ n) -
  of-real (bernoulli (Suc n)) / (real (Suc n) * s ^ n) +
  integral {0..real N} (λx. of-nat n / of-nat (Suc n) *
  (of-real (pbernpoly (Suc n) x) / (of-real x + s) ^ Suc n)))
  → 0 - of-real (bernoulli (Suc n)) / (of-nat (Suc n) * s ^ n) +
  of-nat n / of-nat (Suc n) * stirling-integral (Suc n) s (is ?f →
-)
```

unfolding stirling-integral-def integral-mult-right
using convergent-stirling-integral[of Suc n] **assms** s
by (intro tendsto-intros pbernpoly-over-power-tendsto-0)
(auto simp: convergent-LIMSEQ-iff simp del: of-nat-Suc)
also have ?this ↔ (λN. integral {0..real N}
(λx. of-real (pbernpoly n x) / (of-real x + s) ^ n)) →
of-nat n / of-nat (Suc n) * stirling-integral (Suc n) s -
of-real (bernoulli (Suc n)) / (of-nat (Suc n) * s ^ n)
using eventually-gt-at-top[of 0::nat] pbernpoly-integral-conv-pbernpoly-integral-Suc[of
n]
assms **unfolding** integral-mult-right
by (intro filterlim-cong refl) (auto elim!: eventually-mono simp del: power-Suc)
finally show ?thesis **unfolding** stirling-integral-def[of n] **by** (rule limI)
qed

lemma *stirling-integral-1-unfold*:

assumes $m > 0$

shows $\text{stirling-integral } 1 \ s = \text{stirling-integral } m \ s / \text{of-nat } m -$
 $(\sum_{k=1..<m.} \text{of-real } (\text{bernoulli } (\text{Suc } k)) / (\text{of-nat } k * \text{of-nat } (\text{Suc } k)) *$
 $s \wedge k))$

proof –

have $\text{stirling-integral } 1 \ s = \text{stirling-integral } (\text{Suc } m) \ s / \text{of-nat } (\text{Suc } m) -$
 $(\sum_{k=1..<\text{Suc } m.} \text{of-real } (\text{bernoulli } (\text{Suc } k)) / (\text{of-nat } k * \text{of-nat } (\text{Suc } k))$
 $* s \wedge k))$ **for** m

proof (*induction m*)

case ($\text{Suc } m$)

let $?C = (\sum_{k=1..<\text{Suc } m.} \text{of-real } (\text{bernoulli } (\text{Suc } k)) / (\text{of-nat } k * \text{of-nat}$
 $(\text{Suc } k)) * s \wedge k))$

note *Suc.IH*

also have $\text{stirling-integral } (\text{Suc } m) \ s / \text{of-nat } (\text{Suc } m) =$
 $\text{stirling-integral } (\text{Suc } (\text{Suc } m)) \ s / \text{of-nat } (\text{Suc } (\text{Suc } m)) -$
 $\text{of-real } (\text{bernoulli } (\text{Suc } (\text{Suc } m))) /$
 $(\text{of-nat } (\text{Suc } m) * \text{of-nat } (\text{Suc } (\text{Suc } m))) * s \wedge \text{Suc } m$

(**is** $- = ?A - ?B$) **by** (*subst stirling-integral-conv-stirling-integral-Suc*)
(simp-all del: of-nat-Suc power-Suc add: divide-simps)

also have $?A - ?B - ?C = ?A - (?B + ?C)$ **by** (*rule diff-diff-eq*)

also have $?B + ?C = (\sum_{k=1..<\text{Suc } (\text{Suc } m).} \text{of-real } (\text{bernoulli } (\text{Suc } k)) /$
 $(\text{of-nat } k * \text{of-nat } (\text{Suc } k)) * s \wedge k))$

using s **by** (*simp add: divide-simps*)

finally show $?case$.

qed *simp-all*

note $\text{this}[\text{of } m - 1]$

also from *assms* **have** $\text{Suc } (m - 1) = m$ **by** *simp*

finally show $?thesis$.

qed

lemma *ln-Gamma-stirling-complex*:

assumes $m > 0$

shows $\text{ln-Gamma } s = (s - 1 / 2) * \ln s - s + \ln (2 * \text{pi}) / 2 +$
 $(\sum_{k=1..<m.} \text{of-real } (\text{bernoulli } (\text{Suc } k)) / (\text{of-nat } k * \text{of-nat } (\text{Suc } k)) *$
 $s \wedge k)) -$
 $\text{stirling-integral } m \ s / \text{of-nat } m$

proof –

have $\text{ln-Gamma } s = (s - 1 / 2) * \ln s - s + \ln (2 * \text{pi}) / 2 - \text{stirling-integral}$
 $1 \ s$

using *limI[OF integral-pbernpoly-1]* **by** (*simp add: stirling-integral-def alge-*
bra-simps)

also have $\text{stirling-integral } 1 \ s = \text{stirling-integral } m \ s / \text{of-nat } m -$

$(\sum_{k=1..<m.} \text{of-real } (\text{bernoulli } (\text{Suc } k)) / (\text{of-nat } k * \text{of-nat } (\text{Suc } k))$
 $* s \wedge k))$

using *assms* **by** (*rule stirling-integral-1-unfold*)

finally show $?thesis$ **by** *simp*

qed

lemma *LIMSEQ-stirling-integral*:
 $n > 0 \implies (\lambda x. \text{integral } \{0..real\ } x) (\lambda x. \text{of-real } (pbernpoly\ n\ x) / (\text{of-real } x + s) ^ n))$
 $\longrightarrow \text{stirling-integral } n\ s$ **unfolding** *stirling-integral-def*
using *convergent-stirling-integral[of n]* **by** (*simp only: convergent-LIMSEQ-iff*)

end

lemmas *has-integral-of-real = has-integral-linear[OF - bounded-linear-of-real, unfolded o-def]*
lemmas *integral-of-real = integral-linear[OF - bounded-linear-of-real, unfolded o-def]*

lemma *integrable-ln-Gamma-aux-real*:
assumes $0 < s$
shows $(\lambda x. pbernpoly\ n\ x / (x + s) ^ n)$ *integrable-on* $\{0..real\ } N$
proof –
have $(\lambda x. \text{complex-of-real } (pbernpoly\ n\ x / (x + s) ^ n))$ *integrable-on* $\{0..real\ } N$
using *integrable-ln-Gamma-aux[of of-real s n N]* **assms** **by** *simp*
from *integrable-linear[OF this bounded-linear-Re]* **show** *?thesis*
by (*simp only: o-def Re-complex-of-real*)
qed

lemma
assumes $x > 0\ n > 0$
shows *stirling-integral-complex-of-real*:
 $\text{stirling-integral } n (\text{complex-of-real } x) = \text{of-real } (\text{stirling-integral } n\ x)$
and *LIMSEQ-stirling-integral-real*:
 $(\lambda N. \text{integral } \{0..real\ } N) (\lambda t. pbernpoly\ n\ t / (t + x) ^ n)$
 $\longrightarrow \text{stirling-integral } n\ x$
and *stirling-integral-real-convergent*:
 $\text{convergent } (\lambda N. \text{integral } \{0..real\ } N) (\lambda t. pbernpoly\ n\ t / (t + x) ^ n)$
proof –
have $(\lambda N. \text{integral } \{0..real\ } N) (\lambda t. \text{of-real } (pbernpoly\ n\ t / (t + x) ^ n))$
 $\longrightarrow \text{stirling-integral } n (\text{complex-of-real } x)$
using *LIMSEQ-stirling-integral[of complex-of-real x n]* **assms** **by** *simp*
hence $(\lambda N. \text{of-real } (\text{integral } \{0..real\ } N) (\lambda t. pbernpoly\ n\ t / (t + x) ^ n))$
 $\longrightarrow \text{stirling-integral } n (\text{complex-of-real } x)$
using *integrable-ln-Gamma-aux-real[OF assms(1), of n]*
by (*subst (asm) integral-of-real*) *simp*
from *tendsto-Re[OF this]*
have $(\lambda N. \text{integral } \{0..real\ } N) (\lambda t. pbernpoly\ n\ t / (t + x) ^ n)$
 $\longrightarrow \text{Re } (\text{stirling-integral } n (\text{complex-of-real } x))$ **by** *simp*
thus *convergent* $(\lambda N. \text{integral } \{0..real\ } N) (\lambda t. pbernpoly\ n\ t / (t + x) ^ n)$
by (*rule convergentI*)
thus $(\lambda N. \text{integral } \{0..real\ } N) (\lambda t. pbernpoly\ n\ t / (t + x) ^ n)$
 $\longrightarrow \text{stirling-integral } n\ x$ **unfolding** *stirling-integral-def*
by (*simp add: convergent-LIMSEQ-iff*)

from *tendsto-of-real*[*OF this, where 'a = complex*
integrable-ln-Gamma-aux-real[*OF assms(1), of n*
have ($\lambda x a. \text{integral } \{0..real\ } xa$)
 $(\lambda x a. \text{complex-of-real } (pbernpoly\ n\ xa) / (\text{complex-of-real } xa + x$
 $\wedge n))$
 $\longrightarrow \text{complex-of-real } (\text{stirling-integral } n\ x)$
by (*subst (asm) integral-of-real [symmetric]*) *simp-all*
from *LIMSEQ-unique*[*OF this LIMSEQ-stirling-integral*[*of complex-of-real x n*]]
assms
show *stirling-integral n (complex-of-real x) = of-real (stirling-integral n x)* **by**
simp
qed

lemma *ln-Gamma-stirling-real*:

assumes $x > (0 :: real)$ $m > (0 :: nat)$
shows $\ln\text{-Gamma } x = (x - 1 / 2) * \ln\ x - x + \ln (2 * \pi) / 2 +$
 $(\sum k = 1..<m. \text{bernoulli } (Suc\ k) / (\text{of-nat } k * \text{of-nat } (Suc\ k) * x \wedge k))$
 $-$
 $\text{stirling-integral } m\ x / \text{of-nat } m$

proof $-$

from *assms* **have** *complex-of-real (ln-Gamma x) = ln-Gamma (complex-of-real*
x)
by (*simp add: ln-Gamma-complex-of-real*)
also have $\ln\text{-Gamma } (\text{complex-of-real } x) = \text{complex-of-real } ($
 $(x - 1 / 2) * \ln\ x - x + \ln (2 * \pi) / 2 +$
 $(\sum k = 1..<m. \text{bernoulli } (Suc\ k) / (\text{of-nat } k * \text{of-nat } (Suc\ k) * x \wedge$
 $k)) -$
 $\text{stirling-integral } m\ x / \text{of-nat } m)$ **using** *assms*
by (*subst ln-Gamma-stirling-complex*[*of - m*])
(simp-all add: Ln-of-real stirling-integral-complex-of-real)
finally show *?thesis* **by** (*subst (asm) of-real-eq-iff*)
qed

lemma *stirling-integral-bound-aux*:

assumes $n: n > (1 :: nat)$
obtains c **where** $\bigwedge s. \text{Re } s > 0 \implies \text{norm } (\text{stirling-integral } n\ s) \leq c / \text{Re } s \wedge$
 $(n - 1)$
proof $-$
obtain c **where** $c: \text{norm } (pbernpoly\ n\ x) \leq c$ **for** x **by** (*rule bounded-pbernpoly*[*of*
n]) *blast*
have $c': pbernpoly\ n\ x \leq c$ **for** x **using** c [*of x*] **by** (*simp add: abs-real-def split:*
if-splits)
from c [*of 0*] **have** $c\text{-nonneg}: c \geq 0$ **by** *simp*
have $\text{norm } (\text{stirling-integral } n\ s) \leq c / (\text{real } n - 1) / \text{Re } s \wedge (n - 1)$ **if** $s: \text{Re } s$
 > 0 **for** s
proof (*rule Lim-norm-ubound*[*OF - LIMSEQ-stirling-integral*])
have $\text{pos}: x + \text{norm } s > 0$ **if** $x \geq 0$ **for** x **using** s **that** **by** (*intro add-nonneg-pos*)
auto

have nz : *of-real* $x + s \neq 0$ **if** $x \geq 0$ **for** x **using** s **that by** (*auto simp: complex-eq-iff*)
let $?bound = \lambda N. c / (Re\ s \wedge (n - 1) * (real\ n - 1)) -$
 $c / ((real\ N + Re\ s) \wedge (n - 1) * (real\ n - 1))$
show *eventually* $(\lambda N. norm\ (integral\ \{0..real\ N\}$
 $(\lambda x. of-real\ (pbernpoly\ n\ x) / (of-real\ x + s) \wedge n)) \leq$
 $c / (real\ n - 1) / Re\ s \wedge (n - 1))$ *at-top*
using *eventually-gt-at-top*[*of 0::nat*]
proof *eventually-elim*
case (*elim N*)
let $?F = \lambda x. -c / ((x + Re\ s) \wedge (n - 1) * (real\ n - 1))$
from $n\ s$ **have** $((\lambda x. c / (x + Re\ s) \wedge n)$ *has-integral* $(?F\ (real\ N) - ?F\ 0))$
 $\{0..real\ N\}$
by (*intro fundamental-theorem-of-calculus*)
(auto intro!: derivative-eq-intros simp: divide-simps power-diff add-eq-0-iff2
has-field-derivative-iff-has-vector-derivative [symmetric])
also have $?F\ (real\ N) - ?F\ 0 = ?bound\ N$ **by** *simp*
finally have $*$: $((\lambda x. c / (x + Re\ s) \wedge n)$ *has-integral* $?bound\ N)$ $\{0..real\ N\}$
 \cdot
have $norm\ (integral\ \{0..real\ N\}\ (\lambda x. of-real\ (pbernpoly\ n\ x) / (of-real\ x + s)$
 $\wedge n)) \leq$
 $integral\ \{0..real\ N\}\ (\lambda x. c / (x + Re\ s) \wedge n)$
proof (*intro integral-norm-bound-integral integrable-ln-Gamma-aux s ballI*)
fix x **assume** $x \in \{0..real\ N\}$
have $norm\ (of-real\ (pbernpoly\ n\ x) / (of-real\ x + s) \wedge n) \leq c / norm\ (of-real$
 $x + s) \wedge n$
unfolding *norm-divide norm-power* **using** c **by** (*intro divide-right-mono*)
simp-all
also have $\dots \leq c / (x + Re\ s) \wedge n$
using $x\ c\ c\text{-nonneg}\ s\ nz$ [*of x*] *complex-Re-le-cmod*[*of of-real x + s*]
by (*intro divide-left-mono power-mono mult-pos-pos zero-less-power*
add-nonneg-pos) *auto*
finally show $norm\ (of-real\ (pbernpoly\ n\ x) / (of-real\ x + s) \wedge n) \leq \dots$.
qed (*insert n s * pos nz c, auto simp: complex-nonpos-Reals-iff*)
also have $\dots = ?bound\ N$ **using** $*$ **by** (*simp add: has-integral-iff*)
also have $\dots \leq c / (Re\ s \wedge (n - 1) * (real\ n - 1))$ **using** $c\text{-nonneg}\ elim\ s$
 n **by** *simp*
also have $\dots = c / (real\ n - 1) / (Re\ s \wedge (n - 1))$ **by** *simp*
finally show $norm\ (integral\ \{0..real\ N\}\ (\lambda x. of-real\ (pbernpoly\ n\ x) /$
 $(of-real\ x + s) \wedge n)) \leq c / (real\ n - 1) / Re\ s \wedge (n - 1)$.
qed
qed (*insert s n, simp-all add: complex-nonpos-Reals-iff*)
thus $?thesis$ **by** (*rule that*)
qed

lemma *stirling-integral-bound-aux-integral1*:

fixes $a\ b\ c :: real$ **and** $n :: nat$
assumes $a \geq 0\ b > 0\ c \geq 0\ n > 1\ l < a - b\ r > a + b$
shows $((\lambda x. c / max\ b\ |x - a| \wedge n)$ *has-integral*

$2 * c * (n / (n - 1)) / b \wedge (n - 1) - c / (n - 1) * (1 / (a - l) \wedge (n - 1) + 1 / (r - a) \wedge (n - 1))$
 $\{l..r\}$

proof –

define $x1\ x2$ **where** $x1 = a - b$ **and** $x2 = a + b$

define $F1$ **where** $F1 = (\lambda x :: real. c / (a - x) \wedge (n - 1) / (n - 1))$

define $F3$ **where** $F3 = (\lambda x :: real. -c / (x - a) \wedge (n - 1) / (n - 1))$

have $deriv$: $(F1\ has_vector_derivative\ (c / (a - x) \wedge n))\ (at\ x\ within\ A)$
 $(F3\ has_vector_derivative\ (c / (x - a) \wedge n))\ (at\ x\ within\ A)$

if $x \neq a$ **for** $x :: real$ **and** A

unfolding $F1$ -def $F3$ -def **using** $assms$ **that**

by $(auto\ intro!$: $derivative$ -eq-intros $simp$: $divide$ -simps $power$ -diff add -eq-0-iff2
 $simp\ flip$: has_field -derivative-iff-has-vector-derivative)

from $assms$ **have** $((\lambda x. c / (a - x) \wedge n)\ has_integral\ (F1\ x1 - F1\ l))\ \{l..x1\}$

by $(intro\ fundamental$ -theorem-of-calculus $deriv)$ $(auto\ simp$: $x1$ -def max -def
 $split$: if -splits)

also **have** $?this \longleftrightarrow ((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ (F1\ x1 - F1\ l))$
 $\{l..x1\}$

using $assms$

by $(intro\ has$ -integral-spike-finite-eq[$of\ \{l\}$]) $(auto\ simp$: $x1$ -def max -def $split$:
 if -splits)

finally **have** $I1$: $((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ (F1\ x1 - F1\ l))$
 $\{l..x1\}$.

have $((\lambda x. c / b \wedge n)\ has_integral\ (x2 - x1) * c / b \wedge n)\ \{x1..x2\}$

using has -integral-const-real[$of\ c / b \wedge n\ x1\ x2$] $assms$ **by** $(simp\ add$: $x1$ -def
 $x2$ -def)

also **have** $?this \longleftrightarrow ((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ ((x2 - x1) * c /$
 $b \wedge n))\ \{x1..x2\}$

by $(intro\ has$ -integral-spike-finite-eq[$of\ \{x1, x2\}$])

$(auto\ simp$: $x1$ -def $x2$ -def $split$: if -splits)

finally **have** $I2$: $((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ ((x2 - x1) * c / b \wedge$
 $n))\ \{x1..x2\}$.

from $assms$ **have** $I3$: $((\lambda x. c / (x - a) \wedge n)\ has_integral\ (F3\ r - F3\ x2))\ \{x2..r\}$

by $(intro\ fundamental$ -theorem-of-calculus $deriv)$ $(auto\ simp$: $x2$ -def min -def
 $split$: if -splits)

also **have** $?this \longleftrightarrow ((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ (F3\ r - F3\ x2))$
 $\{x2..r\}$

using $assms$

by $(intro\ has$ -integral-spike-finite-eq[$of\ \{r\}$]) $(auto\ simp$: $x2$ -def min -def $split$:
 if -splits)

finally **have** $I3$: $((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ (F3\ r - F3\ x2))$
 $\{x2..r\}$.

have $((\lambda x. c / max\ b\ |x - a| \wedge n)\ has_integral\ (F1\ x1 - F1\ l) + ((x2 - x1) * c / b \wedge n) + (F3\ r - F3\ x2))\ \{l..r\}$

using $assms$

by $(intro\ has$ -integral-combine[$OF - -\ has$ -integral-combine[$OF - -\ I1\ I2\ I3$])

(auto simp: x1-def x2-def)
also have $(F1\ x1 - F1\ l) + ((x2 - x1) * c / b ^ n) + (F3\ r - F3\ x2) =$
 $F1\ x1 - F1\ l + F3\ r - F3\ x2 + (x2 - x1) * c / b ^ n$
by (simp add: algebra-simps)
also have $x2 - x1 = 2 * b$
using *assms* **by** (simp add: x2-def x1-def min-def max-def)
also have $2 * b * c / b ^ n = 2 * c / b ^ (n - 1)$
using *assms* **by** (simp add: power-diff field-simps)
also have $F1\ x1 - F1\ l + F3\ r - F3\ x2 =$
 $c / (n - 1) * (2 / b ^ (n - 1) - 1 / (a - l) ^ (n - 1) - 1 / (r - a) ^ (n - 1))$
using *assms* **by** (simp add: x1-def x2-def F1-def F3-def field-simps)
also have $\dots + 2 * c / b ^ (n - 1) =$
 $2 * c * (1 + 1 / (n - 1)) / b ^ (n - 1) - c / (n - 1) * (1 / (a - l) ^ (n - 1) +$
 $1 / (r - a) ^ (n - 1))$
using *assms* **by** (simp add: field-simps)
also have $1 + 1 / (n - 1) = n / (n - 1)$
using *assms* **by** (simp add: field-simps)
finally show ?thesis .
qed

lemma *stirling-integral-bound-aux-integral2*:

fixes $a\ b\ c :: \text{real}$ **and** $n :: \text{nat}$

assumes $a \geq 0\ b > 0\ c \geq 0\ n > 1$

obtains I **where** $((\lambda x. c / \max\ b\ |x - a| ^ n) \text{ has-integral } I) \{l..r\}$

$$I \leq 2 * c * (n / (n - 1)) / b ^ (n - 1)$$

proof –

define l' **where** $l' = \min\ l\ (a - b - 1)$

define r' **where** $r' = \max\ r\ (a + b + 1)$

define A **where** $A = 2 * c * (n / (n - 1)) / b ^ (n - 1)$

define B **where** $B = c / \text{real}\ (n - 1) * (1 / (a - l') ^ (n - 1) + 1 / (r' - a) ^ (n - 1))$

have *has-int*: $((\lambda x. c / \max\ b\ |x - a| ^ n) \text{ has-integral } (A - B)) \{l'..r'\}$

using *assms* **unfolding** $A\text{-def}\ B\text{-def}$

by (*intro* *stirling-integral-bound-aux-integral1*) (*auto* *simp*: $l'\text{-def}\ r'\text{-def}$)

have $(\lambda x. c / \max\ b\ |x - a| ^ n) \text{ integrable-on } \{l..r\}$

by (*rule* *integrable-on-subinterval*[*OF* *has-integral-integrable*[*OF* *has-int*]])

(*auto* *simp*: $l'\text{-def}\ r'\text{-def}$)

then obtain I **where** *has-int'*: $((\lambda x. c / \max\ b\ |x - a| ^ n) \text{ has-integral } I) \{l..r\}$

by (*auto* *simp*: *integrable-on-def*)

from *assms* **have** $I \leq A - B$

by (*intro* *has-integral-subset-le*[*OF* - *has-int'* *has-int*]) (*auto* *simp*: $l'\text{-def}\ r'\text{-def}$)

also have $\dots \leq A$

using *assms* **by** (*simp* *add*: $B\text{-def}\ l'\text{-def}\ r'\text{-def}$)

finally show ?thesis **using** *that*[*of* I] *has-int'* **unfolding** $A\text{-def}$ **by** *blast*

qed

```

lemma stirling-integral-bound-aux':
  assumes  $n: n > (1::nat)$  and  $\alpha: \alpha \in \{0 < .. < \pi\}$ 
  obtains  $c$  where  $\bigwedge s::complex. s \in complex-cone' \alpha - \{0\} \implies$ 
     $norm (stirling-integral\ n\ s) \leq c / norm\ s \wedge^{(n-1)}$ 

proof –
  obtain  $c$  where  $c: norm (pbernpoly\ n\ x) \leq c$  for  $x$  by (rule bounded-pbernpoly[of
   $n$ ]) blast
  have  $c'$ :  $pbernpoly\ n\ x \leq c$  for  $x$  using  $c$ [of  $x$ ] by (simp add: abs-real-def split:
  if-splits)
  from  $c$ [of  $0$ ] have  $c\text{-nonneg}: c \geq 0$  by simp

  define  $D$  where  $D = c * Beta (- (real-of-int (- int\ n) / 2) - 1 / 2) (1 / 2)$ 
  / 2
  define  $C$  where  $C = max\ D (2*c*(n/(n-1))/sin\ \alpha \wedge^{(n-1)})$ 

  have  $*$ :  $norm (stirling-integral\ n\ s) \leq C / norm\ s \wedge^{(n-1)}$ 
  if  $s: s \in complex-cone' \alpha - \{0\}$  for  $s::complex$ 
  proof (rule Lim-norm-ubound[OF - LIMSEQ-stirling-integral])
  from  $s\ \alpha$  have  $Arg: |Arg\ s| \leq \alpha$  by (auto simp: complex-cone-altdef)
  have  $s'$ :  $s \notin \mathbb{R}_{\leq 0}$ 
  using complex-cone-inter-nonpos-Reals[of  $-\alpha\ \alpha$ ]  $\alpha\ s$  by auto
  from  $s$  have [simp]:  $s \neq 0$  by auto

  show eventually  $(\lambda N. norm (integral\ \{0..real\ N\}$ 
     $(\lambda x. of-real (pbernpoly\ n\ x) / (of-real\ x + s) \wedge^n)) \leq$ 
     $C / norm\ s \wedge^{(n-1)})$  at-top
  using eventually-gt-at-top[of  $0::nat$ ]
  proof eventually-elim
  case (elim N)
  show ?case
  proof (cases Re s > 0)
  case True
  have int':  $((\lambda x. c * (x^2 + norm\ s^2) powr (-n / 2))\ has-integral$ 
     $D * (norm\ s^2) powr (-n / 2 + 1 / 2))\ \{0<..\}$ 
  using has-integral-mult-left[OF has-integral-Beta3[of  $-n/2\ norm\ s^2$ ],
  of  $c$ ] assms True
  unfolding  $D$ -def by (simp add: algebra-simps)
  hence int':  $((\lambda x. c * (x^2 + norm\ s^2) powr (-n / 2))\ has-integral$ 
     $D * (norm\ s^2) powr (-n / 2 + 1 / 2))\ \{0..\}$ 
  by (subst has-integral-interior [symmetric]) simp-all
  hence integrable:  $(\lambda x. c * (x^2 + norm\ s^2) powr (-n / 2))\ integrable-on$ 
   $\{0..\}$ 
  by (simp add: has-integral-iff)

  have  $norm (integral\ \{0..real\ N\} (\lambda x. of-real (pbernpoly\ n\ x) / (of-real\ x +
  s) \wedge^n)) \leq$ 
     $integral\ \{0..real\ N\} (\lambda x. c * (x^2 + norm\ s^2) powr (-n / 2))$ 
  proof (intro integral-norm-bound-integral s ballI integrable-ln-Gamma-aux)
  have [simp]:  $\{0<..\} - \{0::real..\} = \{\}\ \{0..\} - \{0<..\} = \{0::real\}$ 

```



```

    by auto
    have  $(\lambda x. c * (x^2 + (c \bmod s)^2) \text{ powr } (\text{real-of-int } (- \text{int } n) / 2)) \text{ integrable-on } \{0 < ..\}$ 
      using int by (simp add: has-integral-iff)
      also have  $?this \longleftrightarrow (\lambda x. c * (x^2 + (c \bmod s)^2) \text{ powr } (\text{real-of-int } (- \text{int } n) / 2)) \text{ integrable-on } \{0 ..\}$ 
        by (intro integrable-spike-set-eq) auto
      finally show  $(\lambda x. c * (x^2 + (c \bmod s)^2) \text{ powr } (\text{real-of-int } (- \text{int } n) / 2)) \text{ integrable-on } \{0 .. \text{real } N\}$  by (rule integrable-on-subinterval) auto
  next
  fix x assume  $x \in \{0 .. \text{real } N\}$ 
  have  $\text{nz: complex-of-real } x + s \neq 0$ 
    using True x by (auto simp: complex-eq-iff)
  have  $\text{norm } (\text{of-real } (\text{pbernpoly } n \ x) / (\text{of-real } x + s) ^ n) \leq c / \text{norm } (\text{of-real } x + s) ^ n$ 
    unfolding norm-divide norm-power using c by (intro divide-right-mono)
  simp-all
  also have  $\dots \leq c / \text{sqrt } (x ^ 2 + \text{norm } s ^ 2) ^ n$ 
  proof (intro divide-left-mono mult-pos-pos zero-less-power power-mono)
    show  $\text{sqrt } (x^2 + (c \bmod s)^2) \leq c \bmod (\text{complex-of-real } x + s)$ 
      using x True by (simp add: cmod-def algebra-simps power2-eq-square)
    qed (use x True c-nonneg assms nz in <auto simp: add-nonneg-pos>)
  also have  $\text{sqrt } (x ^ 2 + \text{norm } s ^ 2) ^ n = (x ^ 2 + \text{norm } s ^ 2) \text{ powr } (1/2 * n)$ 
    by (subst powr-powr [symmetric], subst powr-realpow)
      (auto simp: power-half-sqrt add-nonneg-pos)
  also have  $c / \dots = c * (x^2 + \text{norm } s ^ 2) \text{ powr } (-n / 2)$ 
    by (simp add: powr-minus field-simps)
  finally show  $\text{norm } (\text{complex-of-real } (\text{pbernpoly } n \ x) / (\text{complex-of-real } x + s) ^ n) \leq \dots$ 
    qed fact+
  also have  $\dots \leq \text{integral } \{0 ..\} (\lambda x. c * (x^2 + \text{norm } s ^ 2) \text{ powr } (-n / 2))$ 
    using c-nonneg
      by (intro integral-subset-le integrable integrable-on-subinterval[OF integrable]) auto
  also have  $\dots = D * (\text{norm } s ^ 2) \text{ powr } (-n / 2 + 1 / 2)$ 
    using int' by (simp add: has-integral-iff)
  also have  $(\text{norm } s ^ 2) \text{ powr } (-n / 2 + 1 / 2) = \text{norm } s \text{ powr } (2 * (-n / 2 + 1 / 2))$ 
    by (subst powr-powr [symmetric]) auto
  also have  $\dots = \text{norm } s \text{ powr } (-\text{real } (n - 1))$ 
    using assms by (simp add: of-nat-diff)
  also have  $D * \dots = D / \text{norm } s ^ (n - 1)$ 
    by (auto simp: powr-minus powr-realpow field-simps)
  also have  $\dots \leq C / \text{norm } s ^ (n - 1)$ 
    by (intro divide-right-mono) (auto simp: C-def)
  finally show  $\text{norm } (\text{integral } \{0 .. \text{real } N\} (\lambda x. \text{of-real } (\text{pbernpoly } n \ x) / (\text{of-real } x + s) ^ n)) \leq \dots$ 

```

next

case *False*
have $\cos |Arg\ s| = \cos (Arg\ s)$
by (*simp add: abs-if*)
also have $\cos (Arg\ s) = Re (rcis (norm\ s) (Arg\ s)) / norm\ s$
by (*subst Re-rcis*) *auto*
also have $\dots = Re\ s / norm\ s$
by (*subst rcis-cmod-Arg*) *auto*
also have $\dots \leq \cos (pi / 2)$
using *False* **by** (*auto simp: field-simps*)
finally have $|Arg\ s| \geq pi / 2$
using *Arg* α **by** (*subst (asm) cos-mono-le-eq*) *auto*

have $\sin \alpha * norm\ s = \sin (pi - \alpha) * norm\ s$
by *simp*
also have $\dots \leq \sin (pi - |Arg\ s|) * norm\ s$
using α *Arg* $\langle |Arg\ s| \geq pi / 2 \rangle$
by (*intro mult-right-mono sin-monotone-2pi-le*) *auto*
also have $\sin |Arg\ s| \geq 0$
using *Arg-bounded*[*of s*] **by** (*intro sin-ge-zero*) *auto*
hence $\sin (pi - |Arg\ s|) = |\sin |Arg\ s||$
by *simp*
also have $\dots = |\sin (Arg\ s)|$
by (*simp add: abs-if*)
also have $\dots * norm\ s = |Im (rcis (norm\ s) (Arg\ s))|$
by (*simp add: abs-mult*)
also have $\dots = |Im\ s|$
by (*subst rcis-cmod-Arg*) *auto*
finally have *abs-Im-ge*: $|Im\ s| \geq \sin \alpha * norm\ s$.

have [*simp*]: $Im\ s \neq 0 \ s \neq 0$
using $s \langle s \notin \mathbf{R}_{\leq 0} \rangle$ *False*
by (*auto simp: cmod-def zero-le-mult-iff complex-nonpos-Reals-iff*)
have $\sin \alpha > 0$
using *assms* **by** (*intro sin-gt-zero*) *auto*

obtain *I* **where** *I*: $((\lambda x. c / \max |Im\ s| |x + Re\ s| ^ n)$ *has-integral I*)
 $\{0..real\ N\}$

$$I \leq 2 * c * (n / (n - 1)) / |Im\ s| ^ (n - 1)$$

using s *c-nonneg* *assms* *False*

stirling-integral-bound-aux-integral2[*of -Re s |Im s| c n 0 real N*] **by**

auto

have $norm (integral \{0..real\ N\} (\lambda x. of-real (pbernpoly\ n\ x) / (of-real\ x + s) ^ n)) \leq$

$$integral \{0..real\ N\} (\lambda x. c / \max |Im\ s| |x + Re\ s| ^ n)$$

proof (*intro integral-norm-bound-integral integrable-ln-Gamma-aux s ballI*)

```

show ( $\lambda x. c / \max |Im\ s| |x + Re\ s| ^ n$ ) integrable-on {0..real N}
  using I(1) by (simp add: has-integral-iff)
next
  fix x assume x:  $x \in \{0..real\ N\}$ 
  have nz: complex-of-real  $x + s \neq 0$ 
    by (auto simp: complex-eq-iff)
  have norm (complex-of-real (pbernpoly n x) / (complex-of-real  $x + s$ ) ^ n)
    ≤
       $c / \text{norm} (\text{complex-of-real } x + s) ^ n$ 
    unfolding norm-divide norm-power using c[of x] by (intro divide-right-mono) simp-all
    also have  $\dots \leq c / \max |Im\ s| |x + Re\ s| ^ n$ 
      using c-nonneg nz abs-Re-le-cmod[of of-real x + s] abs-Im-le-cmod[of of-real x + s]
      by (intro divide-left-mono power-mono mult-pos-pos zero-less-power)
      (auto simp: less-max-iff-disj)
    finally show norm (complex-of-real (pbernpoly n x) / (complex-of-real  $x + s$ ) ^ n) ≤  $\dots$  .
    qed (auto simp: complex-nonpos-Reals-iff)
    also have  $\dots \leq 2 * c * (n / (n - 1)) / |Im\ s| ^ (n - 1)$ 
      using I by (simp add: has-integral-iff)
    also have  $\dots \leq 2 * c * (n / (n - 1)) / (\sin\ \alpha * \text{norm } s) ^ (n - 1)$ 
      using  $\langle \sin\ \alpha > 0 \rangle$  s c-nonneg abs-Im-ge
      by (intro divide-left-mono mult-pos-pos zero-less-power power-mono mult-nonneg-nonneg) auto
    also have  $\dots = 2 * c * (n / (n - 1)) / \sin\ \alpha ^ (n - 1) / \text{norm } s ^ (n - 1)$ 
      by (simp add: field-simps)
    also have  $\dots \leq C / \text{norm } s ^ (n - 1)$ 
      by (intro divide-right-mono) (auto simp: C-def)
    finally show ?thesis .
  qed
qed
qed (use that assms complex-cone-inter-nonpos-Reals[of - $\alpha$   $\alpha$ ]  $\alpha$  in auto)
thus ?thesis by (rule that)
qed

```

lemma *stirling-integral-bound*:

assumes $n > 0$

obtains *c* **where**

$\bigwedge s. Re\ s > 0 \implies \text{norm} (\text{stirling-integral } n\ s) \leq c / Re\ s ^ n$

proof –

let *?f* = $\lambda s. \text{of-nat } n / \text{of-nat } (Suc\ n) * \text{stirling-integral } (Suc\ n)\ s -$
 $\text{of-real } (\text{bernoulli } (Suc\ n)) / (\text{of-nat } (Suc\ n) * s ^ n)$

from *stirling-integral-bound-aux*[*of Suc n*] **assms** **obtain** *c* **where**

c: $\bigwedge s. Re\ s > 0 \implies \text{norm} (\text{stirling-integral } (Suc\ n)\ s) \leq c / Re\ s ^ n$ **by** *auto*

define *c1* **where** $c1 = \text{real } n / \text{real } (Suc\ n) * c$

define *c2* **where** $c2 = |\text{bernoulli } (Suc\ n)| / \text{real } (Suc\ n)$

have *c2-nonneg*: $c2 \geq 0$ **by** (*simp add: c2-def*)

show *?thesis*

proof (rule that)
fix $s :: \text{complex}$ **assume** $s: \text{Re } s > 0$
hence $s': s \notin \mathbb{R}_{\leq 0}$ **by** (auto simp: complex-nonpos-Reals-iff)
have $\text{stirling-integral } n \ s = ?f \ s$ **using** s' *assms*
by (rule *stirling-integral-conv-stirling-integral-Suc*)
also have $\text{norm } \dots \leq \text{norm } (\text{of-nat } n / \text{of-nat } (\text{Suc } n) * \text{stirling-integral } (\text{Suc } n) \ s) +$
 $\text{norm } (\text{of-real } (\text{bernoulli } (\text{Suc } n)) / (\text{of-nat } (\text{Suc } n) * s ^ n))$
by (rule *norm-triangle-ineq4*)
also have $\dots = \text{real } n / \text{real } (\text{Suc } n) * \text{norm } (\text{stirling-integral } (\text{Suc } n) \ s) +$
 $c2 / \text{norm } s ^ n$ (**is** $- = ?A + ?B$)
by (simp add: *norm-divide norm-mult norm-power c2-def field-simps del: of-nat-Suc*)
also have $?A \leq \text{real } n / \text{real } (\text{Suc } n) * (c / \text{Re } s ^ n)$
by (intro *mult-left-mono c s*) *simp-all*
also have $\dots = c1 / \text{Re } s ^ n$ **by** (simp add: *c1-def*)
also have $c2 / \text{norm } s ^ n \leq c2 / \text{Re } s ^ n$ **using** s *c2-nonneg*
by (intro *divide-left-mono power-mono complex-Re-le-cmod mult-pos-pos zero-less-power*) *auto*
also have $c1 / \text{Re } s ^ n + c2 / \text{Re } s ^ n = (c1 + c2) / \text{Re } s ^ n$
using s **by** (simp add: *field-simps*)
finally show $\text{norm } (\text{stirling-integral } n \ s) \leq (c1 + c2) / \text{Re } s ^ n$ **by** $-$ *simp-all*
qed
qed

lemma *stirling-integral-bound'*:
assumes $n > 0$ **and** $\alpha \in \{0 < \dots < \pi\}$
obtains c **where**
 $\bigwedge s :: \text{complex}. s \in \text{complex-cone}' \ \alpha - \{0\} \implies \text{norm } (\text{stirling-integral } n \ s) \leq c / \text{norm } s ^ n$
proof $-$
let $?f = \lambda s. \text{of-nat } n / \text{of-nat } (\text{Suc } n) * \text{stirling-integral } (\text{Suc } n) \ s -$
 $\text{of-real } (\text{bernoulli } (\text{Suc } n)) / (\text{of-nat } (\text{Suc } n) * s ^ n)$
from *stirling-integral-bound-aux'* [*of Suc n*] *assms* **obtain** c **where**
 $c: \bigwedge s :: \text{complex}. s \in \text{complex-cone}' \ \alpha - \{0\} \implies$
 $\text{norm } (\text{stirling-integral } (\text{Suc } n) \ s) \leq c / \text{norm } s ^ n$ **by** *auto*
define $c1$ **where** $c1 = \text{real } n / \text{real } (\text{Suc } n) * c$
define $c2$ **where** $c2 = |\text{bernoulli } (\text{Suc } n)| / \text{real } (\text{Suc } n)$
have *c2-nonneg*: $c2 \geq 0$ **by** (simp add: *c2-def*)
show *?thesis*
proof (rule that)
fix $s :: \text{complex}$ **assume** $s: s \in \text{complex-cone}' \ \alpha - \{0\}$
have $s': s \notin \mathbb{R}_{\leq 0}$
using *complex-cone-inter-nonpos-Reals* [*of -alpha alpha*] *assms s* **by** *auto*

have $\text{stirling-integral } n \ s = ?f \ s$ **using** s' *assms*
by (intro *stirling-integral-conv-stirling-integral-Suc*) *auto*
also have $\text{norm } \dots \leq \text{norm } (\text{of-nat } n / \text{of-nat } (\text{Suc } n) * \text{stirling-integral } (\text{Suc } n) \ s) +$
 $\text{norm } (\text{of-real } (\text{bernoulli } (\text{Suc } n)) / (\text{of-nat } (\text{Suc } n) * s ^ n))$

$norm (of-real (bernoulli (Suc n)) / (of-nat (Suc n) * s ^ n))$
 by (rule norm-triangle-ineq4)
 also have ... = $real n / real (Suc n) * norm (stirling-integral (Suc n) s) +$
 $c2 / norm s ^ n$ (is - = ?A + ?B)
 by (simp add: norm-divide norm-mult norm-power c2-def field-simps del:
 of-nat-Suc)
 also have ?A ≤ $real n / real (Suc n) * (c / norm s ^ n)$
 by (intro mult-left-mono c s) simp-all
 also have ... = $c1 / norm s ^ n$ by (simp add: c1-def)
 also have $c1 / norm s ^ n + c2 / norm s ^ n = (c1 + c2) / norm s ^ n$
 using s by (simp add: divide-simps)
 finally show $norm (stirling-integral n s) ≤ (c1 + c2) / norm s ^ n$ by -
 simp-all
 qed
 qed

lemma *stirling-integral-holomorphic* [holomorphic-intros]:
 assumes $m: m > 0$ and $A \cap \mathbb{R}_{\leq 0} = \{\}$
 shows *stirling-integral m holomorphic-on A*
proof -
 from *assms* have [simp]: $z \notin \mathbb{R}_{\leq 0}$ if $z \in A$ for z
 using that by auto
 let ?f = $\lambda s::complex. of-nat m * ((s - 1 / 2) * Ln s - s + of-real (ln (2 * pi)$
 $/ 2) +$
 $(\sum k=1..<m. of-real (bernoulli (Suc k)) / (of-nat k * of-nat (Suc k) * s$
 $^ k)) -$
 $ln-Gamma s)$
 have ?f holomorphic-on A using *assms*
 by (auto intro!: holomorphic-intros simp del: of-nat-Suc elim!: nonpos-Reals-cases)
 also have ?this \longleftrightarrow *stirling-integral m holomorphic-on A*
 using *assms* by (intro holomorphic-cong refl)
 (simp-all add: field-simps ln-Gamma-stirling-complex)
 finally show *stirling-integral m holomorphic-on A* .
 qed

lemma *stirling-integral-continuous-on-complex* [continuous-intros]:
 assumes $m: m > 0$ and $A \cap \mathbb{R}_{\leq 0} = \{\}$
 shows *continuous-on A (stirling-integral m :: - \Rightarrow complex)*
 by (intro holomorphic-on-imp-continuous-on *stirling-integral-holomorphic assms*)

lemma *has-field-derivative-stirling-integral-complex*:
 fixes $x :: complex$
 assumes $x \notin \mathbb{R}_{\leq 0}$ $n > 0$
 shows *(stirling-integral n has-field-derivative deriv (stirling-integral n) x) (at*
 $x)$
 using *assms*
 by (intro holomorphic-derivI[OF *stirling-integral-holomorphic*, of $n - \mathbb{R}_{\leq 0}$]) auto

lemma
assumes $n: n > 0$ **and** $x > 0$
shows *deriv-stirling-integral-complex-of-real*:
 $(\text{deriv } \tilde{j}) (\text{stirling-integral } n) (\text{complex-of-real } x) =$
 $\text{complex-of-real } ((\text{deriv } \tilde{j}) (\text{stirling-integral } n) x) \text{ (is ?lhs } x = \text{?rhs } x)$
and *differentiable-stirling-integral-real*:
 $(\text{deriv } \tilde{j}) (\text{stirling-integral } n) \text{ field-differentiable at } x \text{ (is ?thesis2)}$

proof –
let $?A = \{s. \text{Re } s > 0\}$
let $?f = \lambda j x. (\text{deriv } \tilde{j}) (\text{stirling-integral } n) (\text{complex-of-real } x)$
let $?f' = \lambda j x. \text{complex-of-real } ((\text{deriv } \tilde{j}) (\text{stirling-integral } n) x)$

have [*simp*]: *open* $?A$ **by** (*simp add: open-halfspace-Re-gt*)

have $?lhs x = ?rhs x \wedge (\text{deriv } \tilde{j}) (\text{stirling-integral } n) \text{ field-differentiable at } x$
if $x > 0$ **for** x **using** *that*

proof (*induction j arbitrary: x*)
case 0
have $((\lambda x. \text{Re } (\text{stirling-integral } n (\text{of-real } x))) \text{ has-field-derivative}$
 $\text{Re } (\text{deriv } (\lambda x. \text{stirling-integral } n x) (\text{of-real } x))) (\text{at } x) \text{ using } 0 n$
by (*auto intro!: derivative-intros has-vector-derivative-real-field*
field-differentiable-derivI holomorphic-on-imp-differentiable-at[of - ?A]
stirling-integral-holomorphic simp: complex-nonpos-Reals-iff)
also have $?this \longleftrightarrow (\text{stirling-integral } n \text{ has-field-derivative}$
 $\text{Re } (\text{deriv } (\lambda x. \text{stirling-integral } n x) (\text{of-real } x))) (\text{at } x)$
using *eventually-nhds-in-open[of \{0<..\} x] 0 n*
by (*intro has-field-derivative-cong-ev refl*)
(auto elim!: eventually-mono simp: stirling-integral-complex-of-real)
finally have *stirling-integral n field-differentiable at x*
by (*auto simp: field-differentiable-def*)
with $0 n$ **show** $?case$ **by** (*auto simp: stirling-integral-complex-of-real*)

next
case (*Suc j x*)
note $IH = \text{conjunct1}[OF \text{Suc.IH}] \text{conjunct2}[OF \text{Suc.IH}]$
have $*: (\text{deriv } \overset{\sim}{\text{Suc } j}) (\text{stirling-integral } n) (\text{complex-of-real } x) =$
 $\text{of-real } ((\text{deriv } \overset{\sim}{\text{Suc } j}) (\text{stirling-integral } n) x) \text{ if } x: x > 0 \text{ for } x$

proof –
have $\text{deriv } ((\text{deriv } \tilde{j}) (\text{stirling-integral } n)) (\text{complex-of-real } x) =$
 $\text{vector-derivative } (\lambda x. (\text{deriv } \tilde{j}) (\text{stirling-integral } n) (\text{of-real } x)) (\text{at } x)$
using $n x$
by (*intro vector-derivative-of-real-right [symmetric]*
holomorphic-on-imp-differentiable-at[of - ?A] holomorphic-higher-deriv
stirling-integral-holomorphic) (auto simp: complex-nonpos-Reals-iff)
also have $\dots = \text{vector-derivative } (\lambda x. \text{of-real } ((\text{deriv } \tilde{j}) (\text{stirling-integral}$
 $n) x)) (\text{at } x)$
using *eventually-nhds-in-open[of \{0<..\} x] x*
by (*intro vector-derivative-cong-eq) (auto elim!: eventually-mono simp:*

IH(1)
also have ... = of-real (deriv ((deriv \sim j) (stirling-integral n)) x)
by (intro vector-derivative-of-real-left holomorphic-on-imp-differentiable-at[of - ?A]
- ?A]
field-differentiable-imp-differentiable IH(2) x)
finally show ?thesis **by simp**
qed
have ((λx . Re ((deriv \sim Suc j) (stirling-integral n) (of-real x))) has-field-derivative

Re (deriv ((deriv \sim Suc j) (stirling-integral n) (of-real x))) (at x)
using Suc.prem1
by (intro derivative-intros has-vector-derivative-real-field field-differentiable-derivI
holomorphic-on-imp-differentiable-at[of - ?A] stirling-integral-holomorphic
holomorphic-higher-deriv) (auto simp: complex-nonpos-Reals-iff)
also have ?this \longleftrightarrow ((deriv \sim Suc j) (stirling-integral n) has-field-derivative
Re (deriv ((deriv \sim Suc j) (stirling-integral n) (of-real x))) (at x)
using eventually-nhds-in-open[of {0<..} x] Suc.prem1 *
by (intro has-field-derivative-cong-ev refl) (auto elim!: eventually-mono)
finally have (deriv \sim Suc j) (stirling-integral n) field-differentiable at x
by (auto simp: field-differentiable-def)
with *[OF Suc.prem1] **show** ?case **by blast**
qed
from this[OF assms(2)] **show** ?lhs x = ?rhs x ?thesis2 **by blast+**
qed

Unfortunately, asymptotic power series cannot, in general, be differentiated. However, since $\ln\text{-Gamma}$ is holomorphic on the entire positive real half-space, we can differentiate its asymptotic expansion after all.

To do this, we use an ad-hoc version of the more general approach outlined in Erdelyi's "Asymptotic Expansions" for holomorphic functions: We bound the value of the j -th derivative of the remainder term at some value x by applying Cauchy's integral formula along a circle centred at x with radius $\frac{1}{2}x$.

lemma deriv-stirling-integral-real-bound:

assumes $m: m > 0$
shows (deriv \sim j) (stirling-integral m) $\in O(\lambda x::real. 1 / x^{(m+j)})$
proof –
obtain c **where** $c: \bigwedge s. 0 < \text{Re } s \implies c \text{ mod } (\text{stirling-integral } m \ s) \leq c / \text{Re } s^{(m)}$
using stirling-integral-bound[OF m] **by auto**
have $0 \leq c \text{ mod } (\text{stirling-integral } m \ 1)$ **by simp**
also have ... $\leq c$ **using** c[of 1] **by simp**
finally have $c\text{-nonneg}: c \geq 0$.
define B **where** $B = c * 2^{(m + \text{Suc } j)}$
define B' **where** $B' = B * \text{fact } j / 2$

have eventually ($\lambda x::real. \text{norm } ((\text{deriv } \sim j) (\text{stirling-integral } m) x) \leq$
 $B' * \text{norm } (1 / x^{(m+j)})$) at-top

```

using eventually-gt-at-top[of 0::real]
proof eventually-elim
  case (elim x)
  have  $s \notin \mathbb{R}_{\leq 0}$  if  $s \in \text{cball}(\text{of-real } x) (x/2)$  for  $s :: \text{complex}$ 
  proof -
    have  $x - \text{Re } s \leq \text{norm}(\text{of-real } x - s)$  using complex-Re-le-cmod[of of-real x
- s] by simp
    also from that have  $\dots \leq x/2$  by (simp add: dist-complex-def)
    finally show ?thesis using elim by (auto simp: complex-nonpos-Reals-iff)
  qed
  hence  $((\lambda u. \text{stirling-integral } m \ u / (u - \text{of-real } x) \wedge \text{Suc } j)$  has-contour-integral
    complex-of-real  $(2 * \pi) * i / \text{fact } j * (\text{deriv } \hat{\hat{}} j)$  (stirling-integral m) (of-real x)) (circlepath (of-real x) (x/2))
    using m elim
  by (intro Cauchy-has-contour-integral-higher-derivative-circlepath
    stirring-integral-continuous-on-complex stirring-integral-holomorphic)
  auto
  hence  $\text{norm}(\text{of-real } (2 * \pi) * i / \text{fact } j * (\text{deriv } \hat{\hat{}} j)$  (stirling-integral m)
(of-real x))  $\leq$ 
     $B / x \wedge (m + \text{Suc } j) * (2 * \pi * (x / 2))$ 
  proof (rule has-contour-integral-bound-circlepath)
    fix  $u :: \text{complex}$  assume  $\text{dist}(\text{norm}(u - \text{of-real } x) = x / 2)$ 
    have  $\text{Re}(\text{of-real } x - u) \leq \text{norm}(\text{of-real } x - u)$  by (rule complex-Re-le-cmod)
    also have  $\dots = x / 2$  using dist by (simp add: norm-minus-commute)
    finally have  $\text{Re } u: \text{Re } u \geq x/2$  using elim by simp
    have  $\text{norm}(\text{stirling-integral } m \ u / (u - \text{of-real } x) \wedge \text{Suc } j) \leq$ 
       $c / \text{Re } u \wedge m / (x / 2) \wedge \text{Suc } j$  using Re-u elim
    unfolding norm-divide norm-power dist
    by (intro divide-right-mono zero-le-power c) simp-all
    also have  $\dots \leq c / (x/2) \wedge m / (x / 2) \wedge \text{Suc } j$  using c-nonneg elim Re-u
    by (intro divide-right-mono divide-left-mono power-mono) simp-all
    also have  $\dots = B / x \wedge (m + \text{Suc } j)$  using elim by (simp add: B-def
field-simps power-add)
    finally show  $\text{norm}(\text{stirling-integral } m \ u / (u - \text{of-real } x) \wedge \text{Suc } j) \leq B / x$ 
 $\wedge (m + \text{Suc } j)$  .
    qed (insert elim c-nonneg, auto simp: B-def simp del: power-Suc)
    hence  $\text{cmod}((\text{deriv } \hat{\hat{}} j)$  (stirling-integral m) (of-real x))  $\leq B' / x \wedge (j + m)$ 
    using elim by (simp add: field-simps norm-divide norm-mult norm-power
B'-def)
    with elim m show ?case by (simp-all add: add-ac deriv-stirling-integral-complex-of-real)
  qed
  thus ?thesis by (rule bigoI)
qed

```

definition *stirling-sum* **where**

```

  stirling-sum j m x =
     $(-1) \wedge j * (\sum k = 1..<m. (\text{of-real}(\text{bernoulli}(\text{Suc } k)) * \text{pochhammer}(\text{of-nat } k) j / (\text{of-nat } k * \text{of-nat}(\text{Suc } k))) * \text{inverse } x \wedge (k + j))$ 

```


definition *stirling-sum'* where

$$\begin{aligned} \text{stirling-sum}' j m x = \\ (-1) \wedge (\text{Suc } j) * (\sum k \leq m. (\text{of-real } (\text{bernoulli}' k) * \\ \text{pochhammer } (\text{of-nat } (\text{Suc } k)) (j - 1) * \text{inverse } x \wedge (k + j))) \end{aligned}$$

lemma *stirling-sum-complex-of-real*:

$$\begin{aligned} \text{stirling-sum } j m (\text{complex-of-real } x) = \text{complex-of-real } (\text{stirling-sum } j m x) \\ \text{by } (\text{simp add: } \text{stirling-sum-def } \text{pochhammer-of-real } [\text{symmetric}] \text{ del: of-nat-Suc}) \end{aligned}$$

lemma *stirling-sum'-complex-of-real*:

$$\begin{aligned} \text{stirling-sum}' j m (\text{complex-of-real } x) = \text{complex-of-real } (\text{stirling-sum}' j m x) \\ \text{by } (\text{simp add: } \text{stirling-sum}'\text{-def } \text{pochhammer-of-real } [\text{symmetric}] \text{ del: of-nat-Suc}) \end{aligned}$$

lemma *has-field-derivative-stirling-sum-complex* [derivative-intros]:

$\text{Re } x > 0 \implies (\text{stirling-sum } j m \text{ has-field-derivative } \text{stirling-sum } (\text{Suc } j) m x) (\text{at } x)$

$$\begin{aligned} \text{unfolding } \text{stirling-sum-def } [\text{abs-def}] \text{ sum-distrib-left} \\ \text{by } (\text{rule } \text{DERIV-sum}) (\text{auto intro!: } \text{derivative-eq-intros } \text{simp del: of-nat-Suc} \\ \text{simp: } \text{pochhammer-Suc } \text{power-diff}) \end{aligned}$$

lemma *has-field-derivative-stirling-sum-real* [derivative-intros]:

$x > (0::\text{real}) \implies (\text{stirling-sum } j m \text{ has-field-derivative } \text{stirling-sum } (\text{Suc } j) m x) (\text{at } x)$

$$\begin{aligned} \text{unfolding } \text{stirling-sum-def } [\text{abs-def}] \text{ sum-distrib-left} \\ \text{by } (\text{rule } \text{DERIV-sum}) (\text{auto intro!: } \text{derivative-eq-intros } \text{simp del: of-nat-Suc} \\ \text{simp: } \text{pochhammer-Suc } \text{power-diff}) \end{aligned}$$

lemma *has-field-derivative-stirling-sum'-complex* [derivative-intros]:

assumes $j > 0 \text{ Re } x > 0$
shows $(\text{stirling-sum}' j m \text{ has-field-derivative } \text{stirling-sum}' (\text{Suc } j) m x) (\text{at } x)$

proof (cases j)

case $(\text{Suc } j')$

from *assms* **have** [simp]: $x \neq 0$ **by** *auto*

define c **where** $c = (\lambda n. (-1) \wedge \text{Suc } j * \text{complex-of-real } (\text{bernoulli}' n) * \\ \text{pochhammer } (\text{of-nat } (\text{Suc } n)) j')$

define T **where** $T = (\lambda n x. c n * \text{inverse } x \wedge (j + n))$

define T' **where** $T' = (\lambda n x. - (\text{of-nat } (j + n)) * c n * \text{inverse } x \wedge (\text{Suc } (j + n)))$

have $((\lambda x. \sum k \leq m. T k x) \text{ has-field-derivative } (\sum k \leq m. T' k x)) (\text{at } x)$ **using** *assms Suc*

by (*intro DERIV-sum*)

(*auto simp: T-def T'-def intro!: derivative-eq-intros*

simp: field-simps power-add [symmetric] simp del: of-nat-Suc power-Suc

of-nat-add)

also have $(\lambda x. \sum k \leq m. T k x) = \text{stirling-sum}' j m$

by (*simp add: Suc T-def c-def stirling-sum'-def fun-eq-iff add-ac mult.assoc sum-distrib-left*)

also have $(\sum k \leq m. T' k x) = \text{stirling-sum}' (\text{Suc } j) m x$

by (simp add: T'-def c-def Suc stirling-sum'-def sum-distrib-left
sum-distrib-right algebra-simps pochhammer-Suc)
finally show ?thesis .
qed (insert assms, simp-all)

lemma has-field-derivative-stirling-sum'-real [derivative-intros]:

assumes $j > 0$ $x > (0::real)$
shows (stirling-sum' j m has-field-derivative stirling-sum' (Suc j) m x) (at x)
proof (cases j)
case (Suc j')
from assms have [simp]: $x \neq 0$ by auto
define c where $c = (\lambda n. (-1) ^ Suc j * (bernoulli' n) * pochhammer (of-nat (Suc n)) j')$
define T where $T = (\lambda n x. c n * inverse x ^ (j + n))$
define T' where $T' = (\lambda n x. - (of-nat (j + n)) * c n * inverse x ^ (Suc (j + n)))$
have (($\lambda x. \sum k \leq m. T k x$) has-field-derivative ($\sum k \leq m. T' k x$)) (at x) using
assms Suc
by (intro DERIV-sum)
(auto simp: T-def T'-def intro!: derivative-eq-intros
simp: field-simps power-add [symmetric] simp del: of-nat-Suc power-Suc
of-nat-add)
also have ($\lambda x. (\sum k \leq m. T k x) = stirling-sum' j m$)
by (simp add: Suc T-def c-def stirling-sum'-def fun-eq-iff add-ac mult.assoc
sum-distrib-left)
also have ($\sum k \leq m. T' k x = stirling-sum' (Suc j) m x$)
by (simp add: T'-def c-def Suc stirling-sum'-def sum-distrib-left
sum-distrib-right algebra-simps pochhammer-Suc)
finally show ?thesis .
qed (insert assms, simp-all)

lemma higher-deriv-stirling-sum-complex:

$Re\ x > 0 \implies (deriv \widetilde{\sim} i) (stirling-sum\ j\ m)\ x = stirling-sum\ (i + j)\ m\ x$
proof (induction i arbitrary: x)
case (Suc i)
have $deriv ((deriv \widetilde{\sim} i) (stirling-sum\ j\ m))\ x = deriv (stirling-sum\ (i + j)\ m)\ x$
using eventually-nhds-in-open[of {x. Re x > 0} x] Suc.prem
by (intro deriv-cong-ev refl) (auto elim!: eventually-mono simp: open-halfspace-Re-gt
Suc.IH)
also from Suc.prem have ... = $stirling-sum\ (Suc\ (i + j))\ m\ x$
by (intro DERIV-imp-deriv has-field-derivative-stirling-sum-complex)
finally show ?case by simp
qed simp-all

definition Polygamma-approx :: $nat \Rightarrow nat \Rightarrow 'a \Rightarrow 'a :: \{real-normed-field, ln\}$
where

Polygamma-approx j m =
($deriv \widetilde{\sim} j$) ($\lambda x. (x - 1 / 2) * ln\ x - x + of-real (ln (2 * pi)) / 2 +$

stirling-sum 0 m x)

lemma *Polygamma-approx-Suc*: *Polygamma-approx (Suc j) m = deriv (Polygamma-approx j m)*

by (*simp add: Polygamma-approx-def*)

lemma *Polygamma-approx-0*:

*Polygamma-approx 0 m x = (x - 1/2) * ln x - x + of-real (ln (2*pi)) / 2 +*
stirling-sum 0 m x

by (*simp add: Polygamma-approx-def*)

lemma *Polygamma-approx-1-complex*:

Re x > 0 \implies

*Polygamma-approx (Suc 0) m x = ln x - 1 / (2*x) +*
stirling-sum (Suc 0) m x

unfolding *Polygamma-approx-Suc Polygamma-approx-0*

by (*intro DERIV-imp-deriv*)

(*auto intro!: derivative-eq-intros elim!: nonpos-Reals-cases simp: field-simps*)

lemma *Polygamma-approx-1-real*:

x > (0 :: real) \implies

*Polygamma-approx (Suc 0) m x = ln x - 1 / (2*x) +*
stirling-sum (Suc 0) m x

unfolding *Polygamma-approx-Suc Polygamma-approx-0*

by (*intro DERIV-imp-deriv*)

(*auto intro!: derivative-eq-intros elim!: nonpos-Reals-cases simp: field-simps*)

lemma *stirling-sum-2-conv-stirling-sum'-1*:

fixes *x :: 'a :: {real-div-algebra, field-char-0}*

assumes *m > 0 x \neq 0*

shows *stirling-sum' 1 m x = 1 / x + 1 / (2 * x²) +*
stirling-sum 2 m x

proof –

have *pochhammer-2: pochhammer (of-nat k) 2 = of-nat k * of-nat (Suc k)* **for**
k

by (*simp add: pochhammer-Suc eval-nat-numeral add-ac*)

have *stirling-sum 2 m x =*

$(\sum k = \text{Suc } 0..<m. \text{of-real } (\text{bernoulli}' (\text{Suc } k)) * \text{inverse } x^{\text{Suc } (\text{Suc } k)})$

unfolding *stirling-sum-def pochhammer-2 power2-minus power-one mult-1-left*

by (*intro sum.cong refl*)

(*simp-all add: stirling-sum-def pochhammer-2 power2-eq-square divide-simps*
bernoulli'-def

del: of-nat-Suc power-Suc)

also have $1 / (2 * x^2) + \dots =$

$(\sum k=0..<m. \text{of-real } (\text{bernoulli}' (\text{Suc } k)) * \text{inverse } x^{\text{Suc } (\text{Suc } k)})$

using *assms*

by (*subst (2) sum.atLeast-Suc-lessThan*) (*simp-all add: power2-eq-square field-simps*)

also have $1 / x + \dots = (\sum k=0..<\text{Suc } m. \text{of-real } (\text{bernoulli}' k) * \text{inverse } x^{\text{Suc } k})$
Suc k)

by (*subst sum.atLeast0-lessThan-Suc-shift*) (*simp-all add: bernoulli'-def di-*

vide-simps)
also have ... = $(\sum k \leq m. \text{of-real } (\text{bernoulli}' k) * \text{inverse } x \wedge \text{Suc } k)$
by (*intro sum.cong*) *auto*
also have ... = *stirling-sum'* 1 *m x* **by** (*simp add: stirling-sum'-def*)
finally show ?*thesis* **by** (*simp add: add-ac*)
qed

lemma *Polygamma-approx-2-real*:
assumes $x > (0::\text{real})$ $m > 0$
shows *Polygamma-approx* (Suc (Suc 0)) *m x* = *stirling-sum'* 1 *m x*
proof –
have *Polygamma-approx* (Suc (Suc 0)) *m x* = *deriv* (*Polygamma-approx* (Suc 0) *m*) *x*
by (*simp add: Polygamma-approx-Suc*)
also have ... = *deriv* ($\lambda x. \ln x - 1 / (2*x) + \text{stirling-sum } (\text{Suc } 0) \text{ } m \text{ } x$) *x*
using *eventually-nhds-in-open*[of {0<..} *x*] *assms*
by (*intro deriv-cong-ev*) (*auto elim!: eventually-mono simp: Polygamma-approx-1-real*)
also have ... = $1 / x + 1 / (2*x^2) + \text{stirling-sum } (\text{Suc } (\text{Suc } 0)) \text{ } m \text{ } x$ **using**
assms
by (*intro DERIV-imp-deriv*) (*auto intro!: derivative-eq-intros*
elim!: nonpos-Reals-cases simp: field-simps power2-eq-square)
also have ... = *stirling-sum'* 1 *m x* **using** *stirling-sum-2-conv-stirling-sum'-1*[of
m x] *assms*
by (*simp add: eval-nat-numeral*)
finally show ?*thesis* .
qed

lemma *Polygamma-approx-2-complex*:
assumes $\text{Re } x > 0$ $m > 0$
shows *Polygamma-approx* (Suc (Suc 0)) *m x* = *stirling-sum'* 1 *m x*
proof –
have *Polygamma-approx* (Suc (Suc 0)) *m x* = *deriv* (*Polygamma-approx* (Suc 0) *m*) *x*
by (*simp add: Polygamma-approx-Suc*)
also have ... = *deriv* ($\lambda x. \ln x - 1 / (2*x) + \text{stirling-sum } (\text{Suc } 0) \text{ } m \text{ } x$) *x*
using *eventually-nhds-in-open*[of {*s. Re s > 0*} *x*] *assms*
by (*intro deriv-cong-ev*)
(auto simp: open-halfspace-Re-gt elim!: eventually-mono simp: Polygamma-approx-1-complex)
also have ... = $1 / x + 1 / (2*x^2) + \text{stirling-sum } (\text{Suc } (\text{Suc } 0)) \text{ } m \text{ } x$ **using**
assms
by (*intro DERIV-imp-deriv*) (*auto intro!: derivative-eq-intros*
elim!: nonpos-Reals-cases simp: field-simps power2-eq-square)
also have ... = *stirling-sum'* 1 *m x* **using** *stirling-sum-2-conv-stirling-sum'-1*[of
m x] *assms*
by (*subst stirling-sum-2-conv-stirling-sum'-1*) (*auto simp: eval-nat-numeral*)
finally show ?*thesis* .
qed

lemma *Polygamma-approx-ge-2-real*:

```

assumes  $x > (0::\text{real})$   $m > 0$ 
shows  $\text{Polygamma-approx } (\text{Suc } (\text{Suc } j)) \ m \ x = \text{stirling-sum}' (\text{Suc } j) \ m \ x$ 
using  $\text{assms}(1)$ 
proof (induction j arbitrary: x)
  case  $(0 \ x)$ 
    with  $\text{assms}$  show ?case by (simp add: Polygamma-approx-2-real)
  next
    case  $(\text{Suc } j \ x)$ 
      have  $\text{Polygamma-approx } (\text{Suc } (\text{Suc } (\text{Suc } j))) \ m \ x = \text{deriv } (\text{Polygamma-approx } (\text{Suc } (\text{Suc } j)) \ m) \ x$ 
        by (simp add: Polygamma-approx-Suc)
      also have  $\dots = \text{deriv } (\text{stirling-sum}' (\text{Suc } j) \ m) \ x$ 
        using eventually-nhds-in-open[of  $\{0 < ..\}$   $x$ ]  $\text{Suc.prem}$ s
        by (intro deriv-cong-ev refl) (auto elim!: eventually-mono simp: Suc.IH)
      also have  $\dots = \text{stirling-sum}' (\text{Suc } (\text{Suc } j)) \ m \ x$  using  $\text{Suc.prem}$ s
        by (intro DERIV-imp-deriv derivative-intros) simp-all
      finally show ?case .
qed

```

lemma *Polygamma-approx-ge-2-complex:*

```

assumes  $\text{Re } x > 0$   $m > 0$ 
shows  $\text{Polygamma-approx } (\text{Suc } (\text{Suc } j)) \ m \ x = \text{stirling-sum}' (\text{Suc } j) \ m \ x$ 
using  $\text{assms}(1)$ 
proof (induction j arbitrary: x)
  case  $(0 \ x)$ 
    with  $\text{assms}$  show ?case by (simp add: Polygamma-approx-2-complex)
  next
    case  $(\text{Suc } j \ x)$ 
      have  $\text{Polygamma-approx } (\text{Suc } (\text{Suc } (\text{Suc } j))) \ m \ x = \text{deriv } (\text{Polygamma-approx } (\text{Suc } (\text{Suc } j)) \ m) \ x$ 
        by (simp add: Polygamma-approx-Suc)
      also have  $\dots = \text{deriv } (\text{stirling-sum}' (\text{Suc } j) \ m) \ x$ 
        using eventually-nhds-in-open[of  $\{x. \text{Re } x > 0\}$   $x$ ]  $\text{Suc.prem}$ s
        by (intro deriv-cong-ev refl) (auto elim!: eventually-mono simp: Suc.IH open-halfspace-Re-gt)
      also have  $\dots = \text{stirling-sum}' (\text{Suc } (\text{Suc } j)) \ m \ x$  using  $\text{Suc.prem}$ s
        by (intro DERIV-imp-deriv derivative-intros) simp-all
      finally show ?case .
qed

```

lemma *Polygamma-approx-complex-of-real:*

```

assumes  $x > 0$   $m > 0$ 
shows  $\text{Polygamma-approx } j \ m \ (\text{complex-of-real } x) = \text{of-real } (\text{Polygamma-approx } j \ m \ x)$ 
proof (cases j)
  case  $0$ 
    with  $\text{assms}$  show ?thesis by (simp add: Polygamma-approx-0 Ln-of-real stirling-sum-complex-of-real)
  next
    case [simp]:  $(\text{Suc } j')$ 

```

```

thus ?thesis
proof (cases j')
  case 0
    with assms show ?thesis
      by (simp add: Polygamma-approx-1-complex
            Polygamma-approx-1-real stirling-sum-complex-of-real Ln-of-real)
  next
    case (Suc j'')
      with assms show ?thesis
        by (simp add: Polygamma-approx-ge-2-complex Polygamma-approx-ge-2-real
              stirling-sum'-complex-of-real)
qed
qed

lemma higher-deriv-Polygamma-approx [simp]:
  (deriv  $\hat{\sim}$  j) (Polygamma-approx i m) = Polygamma-approx (j + i) m
by (simp add: Polygamma-approx-def funpow-add)

lemma stirling-sum-holomorphic [holomorphic-intros]:
   $0 \notin A \implies$  stirling-sum j m holomorphic-on A
unfolding stirling-sum-def by (intro holomorphic-intros) auto

lemma Polygamma-approx-holomorphic [holomorphic-intros]:
  Polygamma-approx j m holomorphic-on {s. Re s > 0}
unfolding Polygamma-approx-def
by (intro holomorphic-intros) (auto simp: open-halfspace-Re-gt elim!: nonpos-Reals-cases)

lemma higher-deriv-lnGamma-stirling:
  assumes m: m > 0
  shows ( $\lambda x::\text{real}.$  (deriv  $\hat{\sim}$  j) ln-Gamma x - Polygamma-approx j m x)  $\in O(\lambda x.$ 
  1 / x  $^{\wedge}(m + j)$ )
proof -
  have eventually ( $\lambda x.$  |(deriv  $\hat{\sim}$  j) ln-Gamma x - Polygamma-approx j m x| =
    inverse (real m) * |(deriv  $\hat{\sim}$  j) (stirling-integral m) x|) at-top
    using eventually-gt-at-top[of 0::real]
  proof eventually-elim
    case (elim x)
    note x = this
    have  $\forall_F y$  in nhds (complex-of-real x). y  $\in -\mathbb{R}_{\leq 0}$ 
      using elim by (intro eventually-nhds-in-open) auto
    hence (deriv  $\hat{\sim}$  j) ( $\lambda x.$  ln-Gamma x - Polygamma-approx 0 m x) (complex-of-real
  x) =
      (deriv  $\hat{\sim}$  j) ( $\lambda x.$  (-inverse (of-nat m)) * stirling-integral m x)
  (complex-of-real x)
    using x m
    by (intro higher-deriv-cong-ev refl)
      (auto elim!: eventually-mono simp: ln-Gamma-stirling-complex Polygamma-approx-def
        field-simps open-halfspace-Re-gt stirling-sum-def)

```

also have $\dots = - \text{inverse} (\text{of-nat } m) * (\text{deriv } \tilde{j}) (\text{stirling-integral } m) (\text{of-real } x)$ **using** x m
by (*intro higher-deriv-cmult[$\text{of} - \mathbb{R}_{\leq 0}$] *stirling-integral-holomorphic*)
(auto simp: open-halfspace-Re-gt)
also have $(\text{deriv } \tilde{j}) (\lambda x. \text{ln-Gamma } x - \text{Polygamma-approx } 0 \ m \ x) (\text{complex-of-real } x) =$
 $(\text{deriv } \tilde{j}) \text{ln-Gamma} (\text{of-real } x) - (\text{deriv } \tilde{j}) (\text{Polygamma-approx } 0 \ m) (\text{of-real } x)$
using x
by (*intro higher-deriv-diff[$\text{of} - \{s. \text{Re } s > 0\}$] *(auto intro!: holomorphic-intros elim!: nonpos-Reals-cases simp: open-halfspace-Re-gt)*)
also have $(\text{deriv } \tilde{j}) (\text{Polygamma-approx } 0 \ m) (\text{complex-of-real } x) =$
 $\text{of-real} (\text{Polygamma-approx } j \ m \ x)$ **using** x m
by (*simp add: Polygamma-approx-complex-of-real*)
also have $\text{norm} (- \text{inverse} (\text{of-nat } m) * (\text{deriv } \tilde{j}) (\text{stirling-integral } m) (\text{complex-of-real } x)) =$
 $\text{inverse} (\text{real } m) * |(\text{deriv } \tilde{j}) (\text{stirling-integral } m) x|$
using x m **by** (*simp add: norm-mult norm-inverse deriv-stirling-integral-complex-of-real*)
also have $(\text{deriv } \tilde{j}) \text{ln-Gamma} (\text{complex-of-real } x) = \text{of-real} ((\text{deriv } \tilde{j}) \text{ln-Gamma } x)$ **using** x
by (*simp add: higher-deriv-ln-Gamma-complex-of-real*)
also have $\text{norm} (\dots - \text{of-real} (\text{Polygamma-approx } j \ m \ x)) =$
 $|(\text{deriv } \tilde{j}) \text{ln-Gamma } x - \text{Polygamma-approx } j \ m \ x|$
by (*simp only: of-real-diff [symmetric] norm-of-real*)
finally show *?case* .
qed
from *bigthetaI-cong[OF this]* m
have $(\lambda x::\text{real}. (\text{deriv } \tilde{j}) \text{ln-Gamma } x - \text{Polygamma-approx } j \ m \ x) \in$
 $\Theta(\lambda x. (\text{deriv } \tilde{j}) (\text{stirling-integral } m) x)$ **by** *simp*
also have $(\lambda x::\text{real}. (\text{deriv } \tilde{j}) (\text{stirling-integral } m) x) \in O(\lambda x. 1 / x^{(m + j)})$ **using** m
by (*rule deriv-stirling-integral-real-bound*)
finally show *?thesis* .
qed**

lemma *Polygamma-approx-1-real'*:

assumes $x: (\text{real}) > 0$ **and** $m: m > 0$
shows $\text{Polygamma-approx } 1 \ m \ x = \text{ln } x - (\sum k = \text{Suc } 0..m. \text{bernoulli}' \ k * \text{inverse } x^{\wedge} k / \text{real } k)$
proof -
have $\text{Polygamma-approx } 1 \ m \ x = \text{ln } x - (1 / (2 * x) +$
 $(\sum k=\text{Suc } 0..<m. \text{bernoulli} (\text{Suc } k) * \text{inverse } x^{\wedge} \text{Suc } k / \text{real} (\text{Suc } k)))$
(is - = - - (- + ?S)) **using** x **by** (*simp add: Polygamma-approx-1-real *stirling-sum-def**)
also have $?S = (\sum k=\text{Suc } 0..<m. \text{bernoulli}' (\text{Suc } k) * \text{inverse } x^{\wedge} \text{Suc } k / \text{real} (\text{Suc } k))$
by (*intro sum.cong refl*) (*simp-all add: bernoulli'-def*)
also have $1 / (2 * x) + \dots =$
 $(\sum k=0..<m. \text{bernoulli}' (\text{Suc } k) * \text{inverse } x^{\wedge} \text{Suc } k / \text{real} (\text{Suc } k))$

using m
by (*subst* (2) *sum.atLeast-Suc-lessThan*) (*simp-all add: field-simps*)
also have $\dots = (\sum k = \text{Suc } 0..m. \text{bernoulli}' k * \text{inverse } x \wedge k / \text{real } k)$ **using**
assms
by (*subst sum.shift-bounds-Suc-ivl [symmetric]*) (*simp add: atLeastLessThanSuc-atLeastAtMost*)
finally show *?thesis* .
qed

theorem

assumes $m: m > 0$

shows *ln-Gamma-real-asymptotics*:

$(\lambda x. \text{ln-Gamma } x - ((x - 1 / 2) * \text{ln } x - x + \text{ln } (2 * \text{pi}) / 2 +$
 $(\sum k = 1..<m. \text{bernoulli } (\text{Suc } k) / (\text{real } k * \text{real } (\text{Suc } k)) / x \wedge k))$
 $\in O(\lambda x. 1 / x \wedge m)$ (**is** *?th1*)

and *Digamma-real-asymptotics*:

$(\lambda x. \text{Digamma } x - (\text{ln } x - (\sum k=1..m. \text{bernoulli}' k / \text{real } k / x \wedge k))$
 $\in O(\lambda x. 1 / (x \wedge \text{Suc } m))$ (**is** *?th2*)

and *Polygamma-real-asymptotics: $j > 0 \implies$*

$(\lambda x. \text{Polygamma } j x - (-1) \wedge \text{Suc } j * (\sum k \leq m. \text{bernoulli}' k *$
 $\text{pochhammer } (\text{real } (\text{Suc } k)) (j - 1) / x \wedge (k + j))$
 $\in O(\lambda x. 1 / x \wedge (m+j+1))$ (**is** $- \implies$ *?th3*)

proof –

define $G :: \text{nat} \Rightarrow \text{real} \Rightarrow \text{real}$ **where**

$G = (\lambda m. \text{if } m = 0 \text{ then } \text{ln-Gamma} \text{ else } \text{Polygamma } (m - 1))$

have $*$: $(\lambda x. G j x - h x) \in O(\lambda x. 1 / x \wedge (m + j))$

if $\bigwedge x::\text{real}. x > 0 \implies \text{Polygamma-approx } j m x = h x$ **for** $j h$

proof –

have $(\lambda x. G j x - h x) \in$

$\Theta(\lambda x. (\text{deriv } \wedge j) \text{ln-Gamma } x - \text{Polygamma-approx } j m x)$ (**is** $- \in$

$\Theta(?f)$)

using *that*

by (*intro bighetaI-cong*) (*auto intro: eventually-mono[OF eventually-gt-at-top[of 0::real]]*)

simp del: funpow.simps simp: higher-deriv-ln-Gamma-real G-def)

also have $?f \in O(\lambda x::\text{real}. 1 / x \wedge (m + j))$ **using** m

by (*rule higher-deriv-lnGamma-stirling*)

finally show *?thesis* .

qed

note $[[\text{simproc del: simplify-landau-sum}]]$

from $*$ [*OF Polygamma-approx-0*] *assms show ?th1*

by (*simp add: G-def Polygamma-approx-0 stirling-sum-def field-simps*)

from $*$ [*OF Polygamma-approx-1-real'*] *assms show ?th2 by (simp add: G-def field-simps)*

assume $j: j > 0$

from $*$ [*OF Polygamma-approx-ge-2-real, of j - 1*] *assms j show ?th3*

by (*simp add: G-def stirling-sum'-def power-add power-diff field-simps*)

qed

2.5 Asymptotics of the complex Gamma function

The m -th order remainder of Stirling's formula for $\log \Gamma$ is $O(s^{-m})$ uniformly over any complex cone $\text{Arg}(z) \leq \alpha$, $z \neq 0$ for any angle $\alpha \in (0, \pi)$. This means that there is bounded by cz^{-m} for some constant c for all z in this cone.

context

fixes F and α

assumes $\alpha: \alpha \in \{0 < .. < \pi\}$

defines $F \equiv \text{principal}(\text{complex-cone}' \alpha - \{0\})$

begin

lemma *stirling-integral-bigo*:

fixes $m :: \text{nat}$

assumes $m: m > 0$

shows $\text{stirling-integral } m \in O[F](\lambda s. 1 / s \wedge m)$

proof –

obtain c **where** $c: \bigwedge s. s \in \text{complex-cone}' \alpha - \{0\} \implies \text{norm}(\text{stirling-integral } m s) \leq c / \text{norm } s \wedge m$

using *stirling-integral-bound* [*OF* $\langle m > 0 \rangle \alpha$] **by** *blast*

have $0 \leq \text{norm}(\text{stirling-integral } m 1 :: \text{complex})$

by *simp*

also have $\dots \leq c$

using *c[of 1]* α **by** *simp*

finally have $c \geq 0$.

have *eventually* $(\lambda s. s \in \text{complex-cone}' \alpha - \{0\}) F$

unfolding *F-def* **by** (*auto simp: eventually-principal*)

hence *eventually* $(\lambda s. \text{norm}(\text{stirling-integral } m s) \leq c * \text{norm}(1 / s \wedge m)) F$

by *eventually-elim* (*use c in* $\langle \text{simp add: norm-divide norm-power} \rangle$)

thus $\text{stirling-integral } m \in O[F](\lambda s. 1 / s \wedge m)$

by (*intro bigoI[of - c]*) *auto*

qed

end

The following is a more explicit statement of this:

theorem *ln-Gamma-complex-asymptotics-explicit*:

fixes $m :: \text{nat}$ and $\alpha :: \text{real}$

assumes $m > 0$ and $\alpha \in \{0 < .. < \pi\}$

obtains $C :: \text{real}$ and $R :: \text{complex} \Rightarrow \text{complex}$

where $\forall s :: \text{complex}. s \notin \mathbb{R}_{\leq 0} \longrightarrow$

$$\text{ln-Gamma } s = (s - 1/2) * \text{ln } s - s + \text{ln}(2 * \pi) / 2 +$$

$$(\sum k=1..m. \text{bernoulli}(k+1) / (k * (k+1) * s \wedge k)) - R s$$

and $\forall s. s \neq 0 \wedge |\text{Arg } s| \leq \alpha \longrightarrow \text{norm}(R s) \leq C / \text{norm } s \wedge m$

proof –

obtain c **where** $c: \bigwedge s. s \in \text{complex-cone}' \alpha - \{0\} \implies \text{norm}(\text{stirling-integral } m s) \leq c / \text{norm } s \wedge m$

```

    using stirling-integral-bound [OF assms] by blast
  have  $0 \leq \text{norm } (\text{stirling-integral } m \ 1 \ :: \text{ complex})$ 
    by simp
  also have  $\dots \leq c$ 
    using c[of 1] assms by simp
  finally have  $c \geq 0$  .
  define R where  $R = (\lambda s :: \text{ complex. } \text{stirling-integral } m \ s \ / \ \text{of-nat } m)$ 
  show ?thesis
  proof (rule that)
    from ln-Gamma-stirling-complex [of - m] assms show
       $\forall s :: \text{ complex. } s \notin \mathbb{R}_{\leq 0} \longrightarrow$ 
         $\text{ln-Gamma } s = (s - 1 / 2) * \text{ln } s - s + \text{ln } (2 * \text{pi}) / 2 +$ 
         $(\sum_{k=1..<m.} \text{bernoulli } (k+1) / (k * (k+1) * s ^ k)) - R \ s$ 
      by (auto simp add: R-def algebra-simps)
    show  $\forall s. s \neq 0 \wedge |\text{Arg } s| \leq \alpha \longrightarrow \text{cmod } (R \ s) \leq c / \text{real } m / \text{cmod } s ^ m$ 
    proof (safe, goal-cases)
      case (1 s)
      show ?case
        using 1 c[of s] assms
        by (auto simp: complex-cone-altdef abs-le-iff R-def norm-divide field-simps)
    qed
  qed
qed

```

Lastly, we can also derive the asymptotics of Γ itself:

$$\Gamma(z) \sim \sqrt{2\pi/z} \left(\frac{z}{e}\right)^z$$

uniformly for $|z| \rightarrow \infty$ within the cone $\text{Arg}(z) \leq \alpha$ for $\alpha \in (0, \pi)$:

```

context
  fixes F and  $\alpha$ 
  assumes  $\alpha: \alpha \in \{0 < .. < \text{pi}\}$ 
  defines  $F \equiv \text{inf at-infinity } (\text{principal } (\text{complex-cone}' \ \alpha))$ 
begin

```

lemma *Gamma-complex-asymp-equiv:*

Gamma \sim [F] $(\lambda s. \text{sqrt } (2 * \text{pi}) * (s / \text{exp } 1) \text{ powr } s / s \text{ powr } (1 / 2))$

proof –

define *I* :: *complex* \Rightarrow *complex* **where** *I* = *stirling-integral 1*

have *eventually* $(\lambda s. s \in \text{complex-cone}' \ \alpha) \ F$

by (auto *simp: eventually-inf-principal F-def*)

moreover **have** *eventually* $(\lambda s. s \neq 0) \ F$

unfolding *F-def* *eventually-inf-principal*

using *eventually-not-equal-at-infinity* **by** *eventually-elim auto*

ultimately **have** *eventually* $(\lambda s. \text{Gamma } s =$

$\text{sqrt } (2 * \text{pi}) * (s / \text{exp } 1) \text{ powr } s / s \text{ powr } (1 / 2) / \text{exp } (I \ s)) \ F$

proof *eventually-elim*

case (*elim s*)

from *elim* **have** *s'*: $s \notin \mathbb{R}_{\leq 0}$

using *complex-cone-inter-nonpos-Reals*[of $-\alpha$] α **by** *auto*
from *elim* **have** [simp]: $s \neq 0$ **by** *auto*
from s' **have** $\Gamma s = \exp(\ln \Gamma s)$
unfolding *Gamma-complex-altdef* **using** *nonpos-Ints-subset-nonpos-Reals* **by**
auto
also from s' **have** $\ln \Gamma s = (s-1/2) * \text{Ln } s - s + \text{complex-of-real } (\ln$
 $(2 * \pi) / 2) - I s$
by (*subst ln-Gamma-stirling-complex*[of - 1]) (*simp-all add: exp-add exp-diff*
I-def)
also have $\exp \dots = \exp((s-1/2) * \text{Ln } s) / \exp s *$
 $\exp(\text{complex-of-real } (\ln(2 * \pi) / 2)) / \exp(I s)$
unfolding *exp-diff exp-add* **by** (*simp add: exp-diff exp-add*)
also have $\exp((s-1/2) * \text{Ln } s) = s \text{ powr } (s-1/2)$
by (*simp add: powr-def*)
also have $\exp(\text{complex-of-real } (\ln(2 * \pi) / 2)) = \text{sqrt } (2 * \pi)$
by (*subst exp-of-real*) (*auto simp: powr-def simp flip: powr-half-sqrt*)
also have $\exp s = \exp 1 \text{ powr } s$
by (*simp add: powr-def*)
also have $s \text{ powr } (s-1/2) / \exp 1 \text{ powr } s = (s \text{ powr } s / \exp 1 \text{ powr } s) / s$
 $\text{powr } (1/2)$
by (*subst powr-diff*) *auto*
also have $*$: $\text{Ln } (s / \exp 1) = \text{Ln } s - 1$
using *Ln-divide-of-real*[of $\exp 1$ s] **by** (*simp flip: exp-of-real*)
hence $s \text{ powr } s / \exp 1 \text{ powr } s = (s / \exp 1) \text{ powr } s$
unfolding *powr-def* **by** (*subst **) (*auto simp: exp-diff field-simps*)
finally show $\Gamma s = \text{sqrt } (2 * \pi) * (s / \exp 1) \text{ powr } s / s \text{ powr } (1 / 2)$
 $/ \exp(I s)$
by (*simp add: algebra-simps*)
qed
hence $\Gamma s \sim[F] (\lambda s. \text{sqrt } (2 * \pi) * (s / \exp 1) \text{ powr } s / s \text{ powr } (1 / 2) /$
 $\exp(I s))$
by (*rule asymp-equiv-refl-ev*)
also have $\dots \sim[F] (\lambda s. \text{sqrt } (2 * \pi) * (s / \exp 1) \text{ powr } s / s \text{ powr } (1 / 2) / 1)$
proof (*intro asymp-equiv-intros*)
have $F \leq \text{principal } (\text{complex-cone}' \alpha - \{0\})$
unfolding *le-principal F-def eventually-inf-principal*
using *eventually-not-equal-at-infinity* **by** *eventually-elim auto*
moreover have $I \in O[\text{principal } (\text{complex-cone}' \alpha - \{0\})](\lambda s. 1 / s)$
using *stirling-integral-bigo*[of α 1] α **unfolding** *F-def* **by** (*simp add: I-def*)
ultimately have $I \in O[F](\lambda s. 1 / s)$
by (*rule landau-o.big.filter-mono*)
also have $(\lambda s. 1 / s) \in o[F](\lambda s. 1)$
proof (*rule landau-o.smallI*)
fix $c :: \text{real}$
assume $c: c > 0$
hence *eventually* $(\lambda z::\text{complex}. \text{norm } z \geq 1 / c)$ *at-infinity*
by (*auto simp: eventually-at-infinity*)
moreover have *eventually* $(\lambda z::\text{complex}. z \neq 0)$ *at-infinity*
by (*rule eventually-not-equal-at-infinity*)

ultimately show *eventually* $(\lambda z :: \text{complex. norm } (1 / z) \leq c * \text{norm } (1 :: \text{complex})) F$
unfolding *F-def eventually-inf-principal*
by *eventually-elim (use <c > 0> in <auto simp: norm-divide field-simps>)*
qed
finally have $I \in o[F](\lambda s. 1)$.
from *smalloD-tendsto[OF this]* **have** *[tendsto-intros]:* $(I \longrightarrow 0) F$
by *simp*
show $(\lambda x. \text{exp } (I x)) \sim[F] (\lambda x. 1)$
by *(rule asymp-equivI' tendsto-eq-intros refl | simp)+*
qed
finally show *?thesis* **by** *simp*
qed
end
end

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