

# Stirling's formula

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## Abstract

This work contains a proof of Stirling's formula both for the factorial  $n! \sim \sqrt{2\pi n}(n/e)^n$  on natural numbers and the real Gamma function  $\Gamma(x) \sim \sqrt{2\pi/x}(x/e)^x$ . The proof is based on work by Graham Jameson [3].

This is then extended to the full asymptotic expansion

$$\begin{aligned} \log \Gamma(z) = & \left( z - \frac{1}{2} \right) \log z - z + \frac{1}{2} \log(2\pi) + \sum_{k=1}^{n-1} \frac{B_{k+1}}{k(k+1)} z^{-k} \\ & - \frac{1}{n} \int_0^\infty B_n([t])(t+z)^{-n} dt \end{aligned}$$

uniformly for all complex  $z \neq 0$  in the cone  $\arg(z) \leq \alpha$  for any  $\alpha \in (0, \pi)$ , with which the above asymptotic relation for  $\Gamma$  is also extended to complex arguments.

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## 1 Stirling's Formula

```
theory Stirling-Formula
imports
  HOL-Analysis.Analysis
  Landau-Symbols.Landau-More
  HOL-Real-Asymp.Real-Asymp
begin
```

```

context
begin

```

First, we define the  $S'_n$  from Jameson's article:

```

qualified definition  $S' :: nat \Rightarrow real \Rightarrow real$  where

$$S' n x = 1/(2*x) + (\sum r=1..<n. 1/(of-nat r+x)) + 1 / (2*(n + x))$$


```

Next, the trapezium (also called  $T$  in Jameson's article):

```

qualified definition  $T :: real \Rightarrow real$  where

$$T x = 1/(2*x) + 1/(2*(x+1))$$


```

Now we define The difference  $\Delta(x)$ :

```

qualified definition  $D :: real \Rightarrow real$  where

$$D x = T x - ln(x + 1) + ln x$$


```

**qualified lemma**  $S'$ -telescope-trapezium:

```

assumes  $n > 0$ 
shows  $S' n x = (\sum r< n. T (of-nat r + x))$ 
proof (cases n)
  case (Suc m)
    hence  $m : Suc m = n$  by simp
    have  $(\sum r< n. T (of-nat r + x)) =$ 
      
$$(\sum r< Suc m. 1 / (2 * real r + 2 * x)) + (\sum r< n. 1 / (2 * real (Suc r) + 2 * x))$$

    unfolding  $m$  by (simp add: T-def sum.distrib algebra-simps)
    also have  $(\sum r< Suc m. 1 / (2 * real r + 2 * x)) =$ 
      
$$1/(2*x) + (\sum r< m. 1 / (2 * real (Suc r) + 2 * x))$$
 (is  $- = ?a + ?S$ )
    by (subst sum.lessThan-Suc-shift simp)
    also have  $(\sum r< n. 1 / (2 * real (Suc r) + 2 * x)) =$ 
      
$$?S + 1 / (2*(real m + x + 1))$$
 (is  $- = - + ?b$ ) by (simp add: Suc)
    also have  $?a + ?S + (?S + ?b) = 2 * ?S + ?a + ?b$  by (simp add: add-ac)
    also have  $2 * ?S = (\sum r=0..<m. 1 / (real (Suc r) + x))$ 
    unfolding sum-distrib-left by (intro sum.cong auto simp add: divide-simps)
    also have  $(\sum r=0..<m. 1 / (real (Suc r) + x)) = (\sum r=Suc 0..<Suc m. 1 / (real r + x))$ 
    by (subst sum.atLeast-Suc-lessThan-Suc-shift simp-all)
    also have  $\dots = (\sum r=1..<n. 1 / (real r + x))$  unfolding  $m$  by simp
    also have  $\dots + ?a + ?b = S' n x$  by (simp add: S'-def Suc)
    finally show ?thesis ..
qed (insert assms, simp-all)

```

**qualified lemma** *stirling-trapezium*:

```

assumes  $x : (x::real) > 0$ 
shows  $D x \in \{0 .. 1/(12*x^2) - 1/(12 * (x+1)^2)\}$ 
proof -
  define  $y$  where  $y = 1 / (2*x + 1)$ 
  from  $x$  have  $y : y > 0 \wedge y < 1$  by (simp-all add: divide-simps y-def)

```

```

from x have D x = T x - ln ((x + 1) / x) by (subst ln-div) (simp-all add:
D-def)
also from x have (x + 1) / x = 1 + 1 / x by (simp add: field-simps)
finally have D: D x = T x - ln (1 + 1/x) .

from y have ( $\lambda n. y * y^{\wedge}n$ ) sums ( $y * (1 / (1 - y))$ )
  by (intro geometric-sums sums-mult) simp-all
hence ( $\lambda n. y^{\wedge}Suc n$ ) sums ( $y / (1 - y)$ ) by simp
also from x have  $y / (1 - y) = 1 / (2*x)$  by (simp add: y-def divide-simps)
finally have *: ( $\lambda n. y^{\wedge}Suc n$ ) sums ( $1 / (2*x)$ ) .

from y have ( $\lambda n. (-y) * (-y)^{\wedge}n$ ) sums ( $((-y) * (1 / (1 - (-y))))$ )
  by (intro geometric-sums sums-mult) simp-all
hence ( $\lambda n. (-y)^{\wedge}Suc n$ ) sums ( $(-y / (1 + y))$ ) by simp
also from x have  $y / (1 + y) = 1 / (2*(x+1))$  by (simp add: y-def divide-simps)
finally have **: ( $\lambda n. (-y)^{\wedge}Suc n$ ) sums ( $(-1 / (2*(x+1)))$ ) .

from sums-diff[ $OF * **$ ] have sum1: ( $\lambda n. y^{\wedge}Suc n - (-y)^{\wedge}Suc n$ ) sums  $T x$ 
  by (simp add: T-def)

from y have abs y < 1 abs (-y) < 1 by simp-all
from sums-diff[ $OF this[THEN ln-series']$ ]
  have ( $\lambda n. y^{\wedge}n / real n - (-y)^{\wedge}n / real n$ ) sums ( $ln (1 + y) - ln (1 - y)$ )
by simp
also from y have  $ln (1 + y) - ln (1 - y) = ln ((1 + y) / (1 - y))$  by (simp
add: ln-div)
also from x have  $(1 + y) / (1 - y) = 1 + 1/x$  by (simp add: divide-simps
y-def)
finally have ( $\lambda n. y^{\wedge}n / real n - (-y)^{\wedge}n / real n$ ) sums  $ln (1 + 1/x)$  .
hence sum2: ( $\lambda n. y^{\wedge}Suc n / real (Suc n) - (-y)^{\wedge}Suc n / real (Suc n)$ ) sums
 $ln (1 + 1/x)$ 
  by (subst sums-Suc-iff) simp

have  $ln (1 + 1/x) \leq T x$ 
proof (rule sums-le [ $OF - sum2 sum1$ ])
fix n :: nat
show  $y^{\wedge}Suc n / real (Suc n) - (-y)^{\wedge}Suc n / real (Suc n) \leq y^{\wedge}Suc n -$ 
 $(-y)^{\wedge}Suc n$ 
proof (cases even n)
case True
hence eq:  $A - (-y)^{\wedge}Suc n / B = A + y^{\wedge}Suc n / B$   $A - (-y)^{\wedge}Suc n$ 
 $= A + y^{\wedge}Suc n$ 
for A B by simp-all
from y show ?thesis unfolding eq
  by (intro add-mono) (auto simp: divide-simps)
qed simp-all
qed
hence D x ≥ 0 by (simp add: D)

```

```

define c where c = ( $\lambda n.$  if even n then  $2 * (1 - 1 / \text{real}(\text{Suc } n))$  else 0)
note sums-diff[OF sum1 sum2]
also have ( $\lambda n.$   $y \wedge \text{Suc } n - (-y) \wedge \text{Suc } n - (y \wedge \text{Suc } n / \text{real}(\text{Suc } n) -$ 
 $(-y) \wedge \text{Suc } n / \text{real}(\text{Suc } n)) = (\lambda n.$   $c n * y \wedge \text{Suc } n)$ 
by (intro ext) (simp add: c-def algebra-simps)
finally have sum3: ( $\lambda n.$   $c n * y \wedge \text{Suc } n)$  sums D x by (simp add: D)

from y have ( $\lambda n.$   $y \wedge 2 * (\text{of-nat}(\text{Suc } n) * y \wedge n))$  sums ( $y \wedge 2 * (1 / (1 - y) \wedge 2)$ )
by (intro sums-mult geometric-deriv-sums) simp-all
hence ( $\lambda n.$   $\text{of-nat}(\text{Suc } n) * y \wedge (n+2))$  sums ( $y \wedge 2 / (1 - y) \wedge 2$ )
by (simp add: algebra-simps power2-eq-square)
also from x have  $y \wedge 2 / (1 - y) \wedge 2 = 1 / (4 * x \wedge 2)$  by (simp add: y-def divide-simps)
finally have *: ( $\lambda n.$   $\text{real}(\text{Suc } n) * y \wedge (\text{Suc}(\text{Suc } n))$ ) sums ( $1 / (4 * x^2)$ ) by
simp

from y have ( $\lambda n.$   $y \wedge 2 * (\text{of-nat}(\text{Suc } n) * (-y) \wedge n))$  sums ( $y \wedge 2 * (1 / (1 -$ 
 $-y) \wedge 2)$ )
by (intro sums-mult geometric-deriv-sums) simp-all
hence ( $\lambda n.$   $\text{of-nat}(\text{Suc } n) * (-y) \wedge (n+2))$  sums ( $y \wedge 2 / (1 + y) \wedge 2$ )
by (simp add: algebra-simps power2-eq-square)
also from x have  $y \wedge 2 / (1 + y) \wedge 2 = 1 / (2 \wedge 2 * (x+1) \wedge 2)$ 
unfolding power-mult-distrib [symmetric] by (simp add: y-def divide-simps
add-ac)
finally have **: ( $\lambda n.$   $\text{real}(\text{Suc } n) * (-y) \wedge (\text{Suc}(\text{Suc } n))$ ) sums ( $1 / (4 * (x$ 
 $+ 1)^2))$  by simp

define d where d = ( $\lambda n.$  if even n then  $2 * \text{real } n$  else 0)
note sums-diff[OF * **]
also have ( $\lambda n.$   $\text{real}(\text{Suc } n) * y \wedge (\text{Suc}(\text{Suc } n)) - \text{real}(\text{Suc } n) * (-y) \wedge (\text{Suc}(\text{Suc }$ 
 $n))) =$ 
 $(\lambda n.$  d ( $\text{Suc } n) * y \wedge \text{Suc}(\text{Suc } n))$ 
by (intro ext) (simp-all add: d-def)
finally have ( $\lambda n.$   $d n * y \wedge \text{Suc } n)$  sums ( $1 / (4 * x^2) - 1 / (4 * (x + 1)^2)$ )
by (subst (asm) sums-Suc-iff) (simp add: d-def)
from sums-mult[OF this, of 1/3] x
have sum4: ( $\lambda n.$   $d n / 3 * y \wedge \text{Suc } n)$  sums ( $1 / (12 * x \wedge 2) - 1 / (12 * (x$ 
 $+ 1) \wedge 2))$ 
by (simp add: field-simps)

have D x  $\leq (1 / (12 * x \wedge 2) - 1 / (12 * (x + 1) \wedge 2))$ 
proof (intro sums-le [OF - sum3 sum4] allI)
fix n :: nat
define c' :: nat  $\Rightarrow$  real
where c' = ( $\lambda n.$  if odd n  $\vee n = 0$  then 0 else if n = 2 then  $4/3$  else 2)
show c n * y  $\wedge \text{Suc } n \leq d n / 3 * y \wedge \text{Suc } n$ 
proof (intro mult-right-mono)
have c n  $\leq c' n$  by (simp add: c-def c'-def)
also consider n = 0 | n = 1 | n = 2 | n  $\geq 3$  by force

```

```

hence  $c' n \leq d n / 3$  by cases (simp-all add:  $c'$ -def  $d$ -def)
finally show  $c n \leq d n / 3$  .
qed (insert  $y$ , simp)
qed

```

```

with  $\langle D x \geq 0 \rangle$  show ?thesis by simp
qed

```

The following functions correspond to the  $p_n(x)$  from the article. The special case  $n = 0$  would not, strictly speaking, be necessary, but it allows some theorems to work even without the precondition  $n \neq 0$ :

```

qualified definition  $p :: nat \Rightarrow real \Rightarrow real$  where
 $p n x = (if n = 0 then 1/x else (\sum r < n. D (real r + x)))$ 

```

We can write the Digamma function in terms of  $S'$ :

```

qualified lemma  $S'$ -LIMSEQ-Digamma:
assumes  $x: x \neq 0$ 
shows  $(\lambda n. ln (real n) - S' n x - 1/(2*x)) \longrightarrow Digamma x$ 
proof -
define  $c$  where  $c = (\lambda n. ln (real n) - (\sum r < n. inverse (x + real r)))$ 
have eventually  $(\lambda n. 1 / (2 * (x + real n))) = c n - (ln (real n) - S' n x - 1/(2*x))$  at-top
using eventually-gt-at-top[of 0::nat]
proof eventually-elim
fix  $n :: nat$ 
assume  $n: n > 0$ 
have  $c n - (ln (real n) - S' n x - 1/(2*x)) =$ 
 $-(\sum r < n. inverse (real r + x)) + (1/x + (\sum r=1..<n. inverse (real r + x))) + 1/(2*(real n + x))$ 
using  $x$  by (simp add:  $S'$ -def  $c$ -def field-simps)
also have  $1/x + (\sum r=1..<n. inverse (real r + x)) = (\sum r < n. inverse (real r + x))$ 
unfolding lessThan-atLeast0 using  $n$ 
by (subst (2) sum.atLeast-Suc-lessThan) (simp-all add: field-simps)
finally show  $1 / (2 * (x + real n)) = c n - (ln (real n) - S' n x - 1/(2*x))$ 
by simp
qed
moreover have  $(\lambda n. 1 / (2 * (x + real n))) \longrightarrow 0$ 
by real-asymp
ultimately have  $(\lambda n. c n - (ln (real n) - S' n x - 1/(2*x))) \longrightarrow 0$ 
by (blast intro: Lim-transform-eventually)
from tends-to-minus[OF this] have  $(\lambda n. (ln (real n) - S' n x - 1/(2*x)) - c n) \longrightarrow 0$  by simp
moreover from Digamma-LIMSEQ[OF x] have  $c \longrightarrow Digamma x$  by (simp add:  $c$ -def)
ultimately show  $(\lambda n. ln (real n) - S' n x - 1/(2*x)) \longrightarrow Digamma x$ 
by (rule Lim-transform [rotated])
qed

```

Moreover, we can give an expansion of  $S'$  with the  $p$  as variation terms.

**qualified lemma**  $S'$ -approx:

```

 $S' n x = \ln(\text{real } n + x) - \ln x + p n x$ 
proof (cases  $n = 0$ )
  case True
    thus ?thesis by (simp add: p-def  $S'$ -def)
  next
    case False
    hence  $S' n x = (\sum r < n. T(\text{real } r + x))$ 
      by (subst  $S'$ -telescope-trapezium) simp-all
    also have ... =  $(\sum r < n. \ln(\text{real } r + x + 1) - \ln(\text{real } r + x) + D(\text{real } r + x))$ 
      by (simp add: D-def)
    also have ... =  $(\sum r < n. \ln(\text{real } (\text{Suc } r) + x) - \ln(\text{real } r + x)) + p n x$ 
      using False by (simp add: sum.distrib add-ac p-def)
    also have  $(\sum r < n. \ln(\text{real } (\text{Suc } r) + x) - \ln(\text{real } r + x)) = \ln(\text{real } n + x) - \ln x$ 
      by (subst sum-lessThan-telescope) simp-all
    finally show ?thesis .
  qed
```

We define the limit of the  $p$  (simply called  $p(x)$  in Jameson's article):

**qualified definition**  $P :: \text{real} \Rightarrow \text{real}$  **where**

```
 $P x = (\sum n. D(\text{real } n + x))$ 
```

**qualified lemma**  $D$ -summable:

```

assumes  $x: x > 0$ 
shows summable  $(\lambda n. D(\text{real } n + x))$ 
proof -
  have *: summable  $(\lambda n. 1 / (12 * (x + \text{real } n)^2) - 1 / (12 * (x + \text{real } (\text{Suc } n))^2))$ 
    by (rule telescope-summable') real-asymp
  show summable  $(\lambda n. D(\text{real } n + x))$ 
  proof (rule summable-comparison-test[OF - *], rule exI[of - 2], safe)
    fix  $n :: \text{nat}$  assume  $n \geq 2$ 
    show norm  $(D(\text{real } n + x)) \leq 1 / (12 * (x + \text{real } n)^2) - 1 / (12 * (x + \text{real } (\text{Suc } n))^2)$ 
      using stirling-trapezium[of  $\text{real } n + x$ ]  $x$  by (auto simp: algebra-simps)
    qed
  qed
```

**qualified lemma**  $p$ -LIMSEQ:

```

assumes  $x: x > 0$ 
shows  $(\lambda n. p n x) \longrightarrow P x$ 
proof (rule Lim-transform-eventually)
  from  $D$ -summable[OF  $x$ ] have  $(\lambda n. D(\text{real } n + x))$  sums  $P x$  unfolding  $P$ -def
    by (simp add: sums-iff)
  then show  $(\lambda n. \sum r < n. D(\text{real } r + x)) \longrightarrow P x$  by (simp add: sums-def)
  moreover from eventually-gt-at-top[of 1]
```

```

show eventually ( $\lambda n. (\sum r < n. D (\text{real } r + x)) = p n x$ ) at-top
  by eventually-elim (auto simp: p-def)
qed

```

This gives us an expansion of the Digamma function:

```

lemma Digamma-approx:
  assumes  $x : (x :: \text{real}) > 0$ 
  shows Digamma  $x = \ln x - 1 / (2 * x) - P x$ 
proof -
  have eventually ( $\lambda n. \ln(\text{inverse}(1 + x / \text{real } n)) + \ln x - 1 / (2 * x) - p n x =$ 
     $\ln(\text{real } n) - S' n x - 1 / (2 * x)$ ) at-top
  using eventually-gt-at-top[of 1::nat]
  proof eventually-elim
    fix  $n :: \text{nat}$  assume  $n : n > 1$ 
    have  $\ln(\text{real } n) - S' n x = \ln((\text{real } n) / (\text{real } n + x)) + \ln x - p n x$ 
      using assms n unfolding S'-approx by (subst ln-div) (auto simp: algebra_simps)
    also from  $n$  have  $\text{real } n / (\text{real } n + x) = \text{inverse}(1 + x / \text{real } n)$  by (simp add: field-simps)
    finally show  $\ln(\text{inverse}(1 + x / \text{real } n)) + \ln x - 1 / (2 * x) - p n x =$ 
       $\ln(\text{real } n) - S' n x - 1 / (2 * x)$  by simp
  qed
  moreover have ( $\lambda n. \ln(\text{inverse}(1 + x / \text{real } n)) + \ln x - 1 / (2 * x) - p n x$ )
     $\longrightarrow \ln(\text{inverse}(1 + 0)) + \ln x - 1 / (2 * x) - P x$ 
  by (rule tendsto-intros p-LIMSEQ x real-tendsto-divide-at-top
    filterlim-real-sequentially | simp)+
  hence ( $\lambda n. \ln(\text{inverse}(1 + x / \text{real } n)) + \ln x - 1 / (2 * x) - p n x$ )
     $\longrightarrow \ln x - 1 / (2 * x) - P x$  by simp
  ultimately have ( $\lambda n. \ln(\text{real } n) - S' n x - 1 / (2 * x)$ )  $\longrightarrow \ln x - 1 / (2 * x)$ 
   $- P x$ 
  by (blast intro: Lim-transform-eventually)
  moreover from  $x$  have ( $\lambda n. \ln(\text{real } n) - S' n x - 1 / (2 * x)$ )  $\longrightarrow$  Digamma
 $x$ 
  by (intro S'-LIMSEQ-Digamma) simp-all
  ultimately show Digamma  $x = \ln x - 1 / (2 * x) - P x$ 
  by (rule LIMSEQ-unique [rotated])
qed

```

Next, we derive some bounds on  $P$ :

```

qualified lemma p-ge-0:  $x > 0 \implies p n x \geq 0$ 
  using stirling-trapezium[of real n + x for n]
  by (auto simp add: p-def intro!: sum-nonneg)

```

```

qualified lemma P-ge-0:  $x > 0 \implies P x \geq 0$ 
  by (rule tendsto-lowerbound[OF p-LIMSEQ])
  (insert p-ge-0[of x], simp-all)

```

```

qualified lemma p-upper-bound:
  assumes  $x > 0$   $n > 0$ 

```

```

shows  $p n x \leq 1/(12*x^2)$ 
proof –
  from assms have  $p n x = (\sum r < n. D(\text{real } r + x))$ 
    by (simp add: p-def)
  also have ...  $\leq (\sum r < n. 1/(12*(\text{real } r + x)^2) - 1/(12 * (\text{real } (\text{Suc } r) + x)^2))$ 
    using stirling-trapezium[of real r + x for r] assms
    by (intro sum-mono) (simp add: add-ac)
  also have ...  $= 1 / (12 * x^2) - 1 / (12 * (\text{real } n + x)^2)$ 
    by (subst sum-lessThan-telescope') simp
  also have ...  $\leq 1 / (12 * x^2)$  by simp
  finally show ?thesis .
qed

```

**qualified lemma**  $P\text{-upper-bound}$ :

```

assumes  $x > 0$ 
shows  $P x \leq 1/(12*x^2)$ 
proof (rule tendsto-upperbound)
  show eventually  $(\lambda n. p n x \leq 1 / (12 * x^2))$  at-top
    using eventually-gt-at-top[of 0]
    by eventually-elim (use p-upper-bound[of x] assms in auto)
  show  $(\lambda n. p n x) \longrightarrow P x$ 
    by (simp add: assms p-LIMSEQ)
qed auto

```

We can now use this approximation of the Digamma function to approximate *ln-Gamma* (since the former is the derivative of the latter). We therefore define the function  $g$  from Jameson's article, which measures the difference between *ln-Gamma* and its approximation:

```

qualified definition  $g :: \text{real} \Rightarrow \text{real}$  where
   $g x = \text{ln-Gamma } x - (x - 1/2) * \ln x + x$ 

qualified lemma  $\text{DERIV-}g: x > 0 \implies (g \text{ has-field-derivative } -P x) \text{ (at } x)$ 
  unfolding  $g\text{-def [abs-def]}$  using  $\text{Digamma-approx}[of x]$ 
  by (auto intro!: derivative-eq-intros simp: field-simps)

```

**qualified lemma**  $\text{isCont-}P$ :

```

assumes  $x > 0$ 
shows  $\text{isCont } P x$ 
proof –
  define  $D' :: \text{real} \Rightarrow \text{real}$ 
  where  $D' = (\lambda x. -1 / (2 * x^2 * (x+1)^2))$ 
  have  $\text{DERIV-}D: (D \text{ has-field-derivative } D' x) \text{ (at } x)$  if  $x > 0$  for  $x$ 
    unfolding  $D\text{-def [abs-def]} D'\text{-def T-def}$ 
    by (insert that, (rule derivative-eq-intros refl | simp)+)
     $(\text{simp add: power2-eq-square divide-simps, (simp add: algebra-simps) ?})$ 
  note this [THEN DERIV-chain2, derivative-intros]

have  $(P \text{ has-field-derivative } (\sum n. D'(\text{real } n + x))) \text{ (at } x)$ 

```

```

unfolding P-def [abs-def]
proof (rule has-field-derivative-series')
  show convex {x/2<..} by simp
next
  fix n :: nat and y :: real assume y: y ∈ {x/2<..}
  with assms have y > 0 by simp
  thus ((λa. D (real n + a)) has-real-derivative D' (real n + y)) (at y within
  {x/2<..}) by (auto intro!: derivative-eq-intros)
next
  from assms D-summable[of x] show summable (λn. D (real n + x)) by simp
next
  show uniformly-convergent-on {x/2<..} (λn x. ∑ i<n. D' (real i + x))
  proof (rule Weierstrass-m-test')
    fix n :: nat and y :: real
    assume y: y ∈ {x/2<..}
    with assms have y > 0 by auto
    have norm (D' (real n + y)) = (1 / (2 * (y + real n) ^ 2)) * (1 / (y + real
    (Suc n)) ^ 2)
    by (simp add: D'-def add-ac)
    also from y assms have ... ≤ (1 / (2 * (x/2) ^ 2)) * (1 / (real (Suc n)) ^ 2)
    by (intro mult-mono divide-left-mono power-mono) simp-all
    also have 1 / (real (Suc n)) ^ 2 = inverse ((real (Suc n)) ^ 2) by (simp add:
    field-simps)
    finally show norm (D' (real n + y)) ≤ (1 / (2 * (x/2) ^ 2)) * inverse ((real
    (Suc n)) ^ 2) .
next
  show summable (λn. (1 / (2 * (x/2) ^ 2)) * inverse ((real (Suc n)) ^ 2))
  by (subst summable-Suc-iff, intro summable-mult inverse-power-summable)
simp-all
qed
qed (insert assms, simp-all add: interior-open)
thus ?thesis by (rule DERIV-isCont)
qed

qualified lemma P-continuous-on [THEN continuous-on-subset]: continuous-on
{0<..} P
by (intro continuous-at-imp-continuous-on ballII isCont-P) auto

qualified lemma P-integrable:
assumes a: a > 0
shows P integrable-on {a..}
proof -
  define f where f = (λn x. if x ∈ {a..real n} then P x else 0)
  show P integrable-on {a..}
  proof (rule dominated-convergence)
    fix n :: nat
    from a have P integrable-on {a..real n}
    by (intro integrable-continuous-real P-continuous-on) auto

```

```

hence  $f n$  integrable-on  $\{a..real\}$ 
  by (rule integrable-eq) (simp add: f-def)
thus  $f n$  integrable-on  $\{a..\}$ 
  by (rule integrable-on-superset) (auto simp: f-def)
next
  fix  $n :: nat$ 
  show norm ( $f n x$ )  $\leq$  of-real  $(1/12) * (1 / x^2)$  if  $x \in \{a..\}$  for  $x$ 
    using a P-ge-0 P-upper-bound by (auto simp: f-def)
next
  show  $(\lambda x::real. of-real (1/12) * (1 / x^2))$  integrable-on  $\{a..\}$ 
    using has-integral-inverse-power-to-inf[of 2 a] a
    by (intro integrable-on-cmult-left) auto
next
  show  $(\lambda n. f n x) \longrightarrow P x$  if  $x \in \{a..\}$  for  $x$ 
proof -
  have eventually  $(\lambda n. real n \geq x)$  at-top
    using filterlim-real-sequentially by (simp add: filterlim-at-top)
    with that not-frequently have eventually  $(\lambda n. f n x = P x)$  at-top
      by (force simp: f-def)
    thus  $(\lambda n. f n x) \longrightarrow P x$  by (simp add: tends-to-eventually)
  qed
  qed
qed

```

**qualified definition**  $c :: real$  **where**  $c = integral \{1..\} (\lambda x. -P x) + g 1$

We can now give bounds on  $g$ :

```

qualified lemma g-bounds:
assumes  $x: x \geq 1$ 
shows  $g x \in \{c..c + 1/(12*x)\}$ 
proof -
  from assms have int-nonneg: integral  $\{x..\} P \geq 0$ 
    by (intro Henstock-Kurzweil-Integration.integral-nonneg P-integrable)
      (auto simp: P-ge-0)
  have int-upper-bound: integral  $\{x..\} P \leq 1/(12*x)$ 
  proof (rule has-integral-le)
    from x show (P has-integral integral  $\{x..\} P$ )  $\{x..\}$ 
      by (intro integrable-integral P-integrable) simp-all
      from x has-integral-mult-right[OF has-integral-inverse-power-to-inf[of 2 x], of
      1/12]
        show  $((\lambda x. 1/(12*x^2)) has-integral (1/(12*x))) \{x..\}$  by (simp add:
        field-simps)
    qed (insert P-upper-bound x, simp-all)

note DERIV-g [THEN DERIV-chain2, derivative-intros]
from assms have int1:  $((\lambda x. -P x) has-integral (g x - g 1)) \{1..x\}$ 
  by (intro fundamental-theorem-of-calculus)
    (auto simp: has-real-derivative-iff-has-vector-derivative [symmetric]
    intro!: derivative-eq-intros)

```

```

from x have int2: (( $\lambda x. -P x$ ) has-integral integral {x..} ( $\lambda x. -P x$ )) {x..}
  by (intro integrable-integral integrable-neg P-integrable) simp-all
from has-integral-Un[OF int1 int2] x
  have (( $\lambda x. -P x$ ) has-integral g x - g 1 - integral {x..} P) ({1..x}  $\cup$  {x..})
    by (simp add: max-def)
also from x have {1..x}  $\cup$  {x..} = {1..} by auto
finally have (( $\lambda x. -P x$ ) has-integral g x - g 1 - integral {x..} P) {1..} .
moreover have (( $\lambda x. -P x$ ) has-integral integral {1..} ( $\lambda x. -P x$ )) {1..}
  by (intro integrable-integral integrable-neg P-integrable) simp-all
ultimately have g x - g 1 - integral {x..} P = integral {1..} ( $\lambda x. -P x$ )
  by (simp add: has-integral-unique)
hence g x = c + integral {x..} P by (simp add: c-def algebra-simps)
  with int-upper-bound int-nonneg show g x  $\in$  {c..c + 1/(12*x)} by simp
qed

```

Finally, we have bounds on *ln-Gamma*, *Gamma*, and *fact*.

```

qualified lemma ln-Gamma-bounds-aux:
  x  $\geq$  1  $\implies$  ln-Gamma x  $\geq$  c + (x - 1/2) * ln x - x
  x  $\geq$  1  $\implies$  ln-Gamma x  $\leq$  c + (x - 1/2) * ln x - x + 1/(12*x)
  using g-bounds[of x] by (simp-all add: g-def)

qualified lemma Gamma-bounds-aux:
  assumes x: x  $\geq$  1
  shows Gamma x  $\geq$  exp c * x powr (x - 1/2) / exp x
    Gamma x  $\leq$  exp c * x powr (x - 1/2) / exp x * exp (1/(12*x))
proof -
  have exp (ln-Gamma x)  $\geq$  exp (c + (x - 1/2) * ln x - x)
    by (subst exp-le-cancel-iff, rule ln-Gamma-bounds-aux) (simp add: x)
  with x show Gamma x  $\geq$  exp c * x powr (x - 1/2) / exp x
    by (simp add: Gamma-real-pos-exp exp-add exp-diff powr-def del: exp-le-cancel-iff)
next
  have exp (ln-Gamma x)  $\leq$  exp (c + (x - 1/2) * ln x - x + 1/(12*x))
    by (subst exp-le-cancel-iff, rule ln-Gamma-bounds-aux) (simp add: x)
  with x show Gamma x  $\leq$  exp c * x powr (x - 1/2) / exp x * exp (1/(12*x))
    by (simp add: Gamma-real-pos-exp exp-add exp-diff powr-def del: exp-le-cancel-iff)
qed

```

```

qualified lemma Gamma-asymp-equiv-aux:
  Gamma ~[at-top] ( $\lambda x. \exp c * x \text{ powr } (x - 1/2) / \exp x$ )
proof (rule asymp-equiv-sandwich)
  include asymp-equiv-syntax
  show eventually ( $\lambda x. \exp c * x \text{ powr } (x - 1/2) / \exp x \leq \text{Gamma } x$ ) at-top
    eventually ( $\lambda x. \exp c * x \text{ powr } (x - 1/2) / \exp x * \exp (1/(12*x)) \geq \text{Gamma } x$ )
  at-top
  using eventually-ge-at-top[of 1::real]
  by (eventually-elim; use Gamma-bounds-aux in force)+
  have (( $\lambda x::\text{real}. \exp (1 / (12 * x))$ ) —> exp 0) at-top
    by real-asymp
  hence ( $\lambda x. \exp (1 / (12 * x))$ ) ~ ( $\lambda x. 1 :: \text{real}$ )

```

```

    by (intro asymp-equivI') simp-all
  hence ( $\lambda x. \exp c * x \text{powr} (x - 1 / 2) / \exp x * 1$ ) ~
        ( $\lambda x. \exp c * x \text{powr} (x - 1 / 2) / \exp x * \exp (1 / (12 * x))$ )
    by (intro asymp-equiv-mult asymp-equiv-refl) (simp add: asymp-equiv-sym)
  thus ( $\lambda x. \exp c * x \text{powr} (x - 1 / 2) / \exp x$ ) ~
        ( $\lambda x. \exp c * x \text{powr} (x - 1 / 2) / \exp x * \exp (1 / (12 * x))$ ) by simp
qed simp-all

qualified lemma exp-1-powr-real [simp]:  $\exp (1::\text{real}) \text{powr} x = \exp x$ 
  by (simp add: powr-def)

qualified lemma fact-asymp-equiv-aux:
  fact ~[at-top] ( $\lambda x. \exp c * \text{sqrt} (\text{real } x) * (\text{real } x / \exp 1) \text{powr real } x$ )
proof -
  include asymp-equiv-syntax
  have fact ~( $\lambda n. \text{Gamma} (\text{real} (\text{Suc } n))$ ) by (simp add: Gamma-fact)
  also have eventually ( $\lambda n. \text{Gamma} (\text{real} (\text{Suc } n)) = \text{real } n * \text{Gamma} (\text{real } n)$ )
at-top
  using eventually-gt-at-top[of 0::nat]
  by eventually-elim (insert Gamma-plus1[of real n for n],
    auto simp: add-ac of-nat-in-nonpos-Ints-iff)
  also have ( $\lambda n. \text{Gamma} (\text{real } n)$ ) ~ ( $\lambda n. \exp c * \text{real } n \text{powr} (\text{real } n - 1/2) / \exp (\text{real } n)$ )
    by (rule asymp-equiv-compose'[OF Gamma-asymp-equiv-aux] filterlim-real-sequentially)+
  also have eventually ( $\lambda n. \text{real } n * (\exp c * \text{real } n \text{powr} (\text{real } n - 1 / 2) / \exp (\text{real } n)) = \exp c * \text{sqrt} (\text{real } n) * (\text{real } n / \exp 1) \text{powr real } n$ ) at-top
  using eventually-gt-at-top[of 0::nat]
proof eventually-elim
  fix n :: nat assume n:  $n > 0$ 
  thus  $\text{real } n * (\exp c * \text{real } n \text{powr} (\text{real } n - 1 / 2) / \exp (\text{real } n)) = \exp c * \text{sqrt} (\text{real } n) * (\text{real } n / \exp 1) \text{powr real } n$ 
    by (subst powr-diff) (simp-all add: powr-divide powr-half-sqrt field-simps)
qed
finally show ?thesis by - (simp-all add: asymp-equiv-mult)
qed

```

We can also bound *Digamma* above and below.

```

lemma Digamma-plus-1-gt-ln:
  assumes x:  $x > (0 :: \text{real})$ 
  shows  $\text{Digamma} (x + 1) > \ln x$ 
proof -
  have 0 < (17 :: real)
    by simp
  also have 17 ≤ 12 * x ^ 2 + 28 * x + 17
    using x by auto
  finally have 0 < (12 * x ^ 2 + 28 * x + 17) / (12 * (x + 1) ^ 2 * (1 + 2 * x))
    using x by (intro divide-pos-pos mult-pos-pos zero-less-power) auto

```

```

also have ... =  $2 / (2 * x + 1) - 1 / (2 * (x + 1)) - 1 / (12 * (x + 1) ^ 2)$ 
  using x by (simp add: divide-simps) (auto simp: field-simps power2-eq-square
add-eq-0-iff)
also have  $2 / (2 * x + 1) \leq \ln(x + 1) - \ln x$ 
  using ln-inverse-approx-ge[of x x + 1] x by simp
also have ... =  $1 / (2 * (x + 1)) - 1 / (12 * (x + 1) ^ 2) \leq$ 
   $\ln(x + 1) - \ln x - 1 / (2 * (x + 1)) - P(x + 1)$ 
  using P-upper-bound[of x + 1] x by (intro diff-mono) auto
also have ... = Digamma(x + 1) - ln x
  by (subst Digamma-approx) (use x in auto)
finally show Digamma(x + 1) > ln x
  by simp
qed

```

**lemma** Digamma-less-ln:

```

assumes x:  $x > (0 :: real)$ 
shows Digamma x < ln x
proof -
  have Digamma x - ln x =  $- (1 / (2 * x)) - P x$ 
    by (subst Digamma-approx) (use x in auto)
  also have ... < 0 - P x
    using x by (intro diff-strict-right-mono) auto
  also have ... ≤ 0
    using P-ge-0[of x] x by simp
  finally show Digamma x < ln x
    by simp
qed

```

We still need to determine the constant term  $c$ , which we do using Wallis' product formula:

$$\prod_{n=1}^{\infty} \frac{4n^2}{4n^2 - 1} = \frac{\pi}{2}$$

**qualified lemma** powr-mult-2:  $(x::real) > 0 \implies x \text{ powr } (y * 2) = (x ^ 2) \text{ powr } y$ 
**by** (subst mult.commute, subst powr-powr [symmetric]) (simp add: powr-numeral)

**qualified lemma** exp-mult-2:  $\exp(y * 2 :: real) = \exp y * \exp y$ 
**by** (subst exp-add [symmetric]) simp

**qualified lemma** exp-c:  $\exp c = \sqrt{2*pi}$

```

proof -
  include asymp-equiv-syntax
  define p where p = ( $\lambda n. \prod_{k=1..n} (4 * real k ^ 2) / (4 * real k ^ 2 - 1)$ )

```

```

have p-0 [simp]: p 0 = 1 by (simp add: p-def)
have p-Suc: p (Suc n) = p n * (4 * real (Suc n) ^ 2) / (4 * real (Suc n) ^ 2 - 1)
  for n unfolding p-def by (subst prod.nat-ivl-Suc') simp-all
have p: p = ( $\lambda n. 16 ^ n * fact n ^ 4 / (fact(2 * n))^2 / (2 * real n + 1)$ )
proof

```

```

fix n :: nat
have p n = ( $\prod_{k=1..n} (2 * \text{real } k)^2 / (2 * \text{real } k - 1) / (2 * \text{real } k + 1)$ )
  unfolding p-def by (intro prod.cong refl) (simp add: field-simps power2-eq-square)
  also have ... = ( $\prod_{k=1..n} (2 * \text{real } k)^2 / (2 * \text{real } k - 1) / (\prod_{k=1..n} (2 * \text{real } (\text{Suc } k) - 1))$ )
    by (simp add: prod-divide prod.distrib add-ac)
  also have ( $\prod_{k=1..n} (2 * \text{real } (\text{Suc } k) - 1) = (\prod_{k=\text{Suc } 1..n} (2 * \text{real } k - 1))$ )
    by (subst prod.atLeast-Suc-atMost-Suc-shift) simp-all
  also have ... = ( $\prod_{k=1..n} (2 * \text{real } k - 1) * (2 * \text{real } n + 1)$ )
    by (induction n) (simp-all add: prod.nat-ivl-Suc')
  also have ( $\prod_{k=1..n} (2 * \text{real } k)^2 / (2 * \text{real } k - 1) / \dots =$ 
 $(\prod_{k=1..n} (2 * \text{real } k)^2 / (2 * \text{real } k - 1)^2) / (2 * \text{real } n + 1)$ )
  unfolding power2-eq-square by (simp add: prod.distrib prod-divide)
  also have ( $\prod_{k=1..n} (2 * \text{real } k)^2 / (2 * \text{real } k - 1)^2 =$ 
 $(\prod_{k=1..n} (2 * \text{real } k)^4 / ((2 * \text{real } k) * (2 * \text{real } k - 1))^2)$ )
    by (rule prod.cong) (simp-all add: power2-eq-square eval-nat-numeral)
  also have ... =  $16^{\wedge}n * \text{fact } n^{\wedge}4 / (\prod_{k=1..n} (2 * \text{real } k) * (2 * \text{real } k - 1))^2$ 
    by (simp add: prod.distrib prod-divide fact-prod
      prod-power-distrib [symmetric] prod-constant)
  also have ( $\prod_{k=1..n} (2 * \text{real } k) * (2 * \text{real } k - 1) = \text{fact } (2 * n)$ )
    by (induction n) (simp-all add: algebra-simps prod.nat-ivl-Suc')
  finally show p n =  $16^{\wedge}n * \text{fact } n^{\wedge}4 / (\text{fact } (2 * n))^2 / (2 * \text{real } n + 1)$ .
qed

have p ~ ( $\lambda n. 16^{\wedge}n * \text{fact } n^{\wedge}4 / (\text{fact } (2 * n))^2 / (2 * \text{real } n + 1)$ )
  by (simp add: p)
also have ... ~ ( $\lambda n. 16^{\wedge}n * (\exp c * \sqrt{\text{real } n}) * (\text{real } n / \exp 1) \text{powr real } n^{\wedge}4 /$ 
 $(\exp c * \sqrt{\text{real } (2 * n)}) * (\text{real } (2 * n) / \exp 1) \text{powr real } (2 * n)^{\wedge}2 /$ 
 $(2 * \text{real } n + 1)$ ) (is - ~ ?f)
by (intro asymp-equiv-mult asymp-equiv-divide asymp-equiv-refl mult-nat-left-at-top
  fact-asymp-equiv-aux asymp-equiv-power asymp-equiv-compose'[OF
  fact-asymp-equiv-aux])
  simp-all
also have eventually ( $\lambda n. \dots n = \exp c^{\wedge}2 / (4 + 2/n)$ ) at-top
  using eventually-gt-at-top[of 0::nat]
proof eventually-elim
fix n :: nat assume n:  $n > 0$ 
have [simp]:  $16^{\wedge}n = 4^{\wedge}n * (4^{\wedge}n :: \text{real})$  by (simp add: power-mult-distrib
[symmetric])
from n have ?f n =  $\exp c^{\wedge}2 * (n / (2 * (2 * n + 1)))$ 
  by (simp add: power-mult-distrib divide-simps powr-mult real-sqrt-power-even)
    (simp add: field-simps power2-eq-square eval-nat-numeral powr-mult-2
      exp-mult-2 powr-realpow)
also from n have ... =  $\exp c^{\wedge}2 / (4 + 2/n)$  by (simp add: field-simps)
finally show ?f n = ... .
qed

```

```

also have ( $\lambda x. 4 + 2 / \text{real } x$ )  $\sim$  ( $\lambda x. 4$ )
  by (subst asymp-equiv-add-right) auto
finally have  $p \longrightarrow \exp c \wedge 2 / 4$ 
  by (rule asymp-equivD-const) (simp-all add: asymp-equiv-divide)
moreover have  $p \longrightarrow pi / 2$  unfolding p-def by (rule wallis)
ultimately have  $\exp c \wedge 2 / 4 = pi / 2$  by (rule LIMSEQ-unique)
hence  $2 * pi = \exp c \wedge 2$  by simp
also have  $\sqrt{\exp c \wedge 2} = \exp c$  by simp
finally show  $\exp c = \sqrt{2 * pi}$  ..
qed

```

**qualified lemma**  $c: c = \ln(2*pi) / 2$

```

proof -
  note exp-c [symmetric]
  also have  $\ln(\exp c) = c$  by simp
  finally show ?thesis by (simp add: ln-sqrt)
qed

```

This gives us the final bounds:

**theorem** Gamma-bounds:

```

assumes  $x \geq 1$ 
shows  $\Gamma(x) \geq \sqrt{2*pi/x} * (x / \exp 1)$  powr x (is ?th1)
         $\Gamma(x) \leq \sqrt{2*pi/x} * (x / \exp 1)$  powr x * exp(1 / (12 * x)) (is
?th2)
proof -
  from assms have  $\exp c * x \text{powr}(x - 1/2) / \exp x = \sqrt{2*pi/x} * (x / \exp$ 
 $1) \text{powr } x$ 
    by (subst powr-diff)
      (simp add: exp-c real-sqrt-divide powr-divide powr-half-sqrt field-simps)
  with Gamma-bounds-aux[OF assms] show ?th1 ?th2 by simp-all
qed

```

**theorem** ln-Gamma-bounds:

```

assumes  $x \geq 1$ 
shows  $\ln\Gamma(x) \geq \ln(2*pi/x) / 2 + x * \ln x - x$  (is ?th1)
         $\ln\Gamma(x) \leq \ln(2*pi/x) / 2 + x * \ln x - x + 1/(12*x)$  (is ?th2)
proof -
  from ln-Gamma-bounds-aux[OF assms] assms show ?th1 ?th2
    by (simp-all add: c field-simps ln-div)
  from assms have  $\exp c * x \text{powr}(x - 1/2) / \exp x = \sqrt{2*pi/x} * (x / \exp$ 
 $1) \text{powr } x$ 
    by (subst powr-diff)
      (simp add: exp-c real-sqrt-divide powr-divide powr-half-sqrt field-simps)
qed

```

**theorem** fact-bounds:

```

assumes  $n > 0$ 
shows  $(\text{fact } n :: \text{real}) \geq \sqrt{2*pi*n} * (n / \exp 1) \wedge n$  (is ?th1)
         $(\text{fact } n :: \text{real}) \leq \sqrt{2*pi*n} * (n / \exp 1) \wedge n * \exp(1 / (12 * n))$  (is
?th2)

```

```

?th2)
proof -
  from assms have n: real n ≥ 1 by simp
  from assms Gamma-plus1[of real n]
    have real n * Gamma (real n) = Gamma (real (Suc n))
      by (simp add: of-nat-in-nonpos-Ints-iff add-ac)
    also have Gamma (real (Suc n)) = fact n by (subst Gamma-fact [symmetric])
  simp
    finally have *: fact n = real n * Gamma (real n) by simp

  have 2*pi/n = 2*pi*n / n ^ 2 by (simp add: power2-eq-square)
  also have sqrt ... = sqrt (2*pi*n) / n by (subst real-sqrt-divide) simp-all
  also have real n * ... = sqrt (2*pi*n) by simp
  finally have **: real n * sqrt (2*pi/real n) = sqrt (2*pi*real n) .

  note *
  also note Gamma-bounds(2)[OF n]
  also have real n * (sqrt (2 * pi / real n) * (real n / exp 1) powr real n *
    exp (1 / (12 * real n))) =
    (real n * sqrt (2*pi/n)) * (n / exp 1) powr n * exp (1 / (12 * n))
    by (simp add: algebra-simps)
  also from n have (real n / exp 1) powr real n = (real n / exp 1) ^ n
    by (subst powr-realpow) simp-all
  also note **
  finally show ?th2 by - (insert assms, simp-all)

  have sqrt (2*pi*n) * (n / exp 1) powr n = n * (sqrt (2*pi/n) * (n / exp 1)
  powr n)
    by (subst ** [symmetric]) (simp add: field-simps)
  also from assms have ... ≤ real n * Gamma (real n)
    by (intro mult-left-mono Gamma-bounds(1)) simp-all
  also from n have (real n / exp 1) powr real n = (real n / exp 1) ^ n
    by (subst powr-realpow) simp-all
  also note * [symmetric]
  finally show ?th1 .

qed

theorem ln-fact-bounds:
  assumes n > 0
  shows ln (fact n :: real) ≥ ln (2*pi*n)/2 + n * ln n - n (is ?th1)
    ln (fact n :: real) ≤ ln (2*pi*n)/2 + n * ln n - n + 1/(12*real n) (is
?th2)
proof -
  have ln (fact n :: real) ≥ ln (sqrt (2*pi*real n)*(real n/exp 1)^n)
    using fact-bounds(1)[OF assms] assms by (subst ln-le-cancel-iff) auto
  also from assms have ln (sqrt (2*pi*real n)*(real n/exp 1)^n) = ln (2*pi*n)/2
+ n * ln n - n
    by (simp add: ln-mult ln-div ln-realpow ln-sqrt field-simps)
  finally show ?th1 .

```

```

next
  have  $\ln(\text{fact } n :: \text{real}) \leq \ln(\sqrt(2*pi*\text{real } n) * (\text{real } n/\exp 1)^n * \exp(1/(12*\text{real } n)))$ 
    using fact-bounds(2)[OF assms] assms by (subst ln-le-cancel-iff) auto
  also from assms have ... =  $\ln(2*pi*n)/2 + n * \ln n - n + 1/(12*\text{real } n)$ 
    by (simp add: ln-mult ln-div ln-realpow ln-sqrt field-simps)
  finally show ?th2 .
qed

```

```

theorem Gamma-asymp-equiv:
   $\Gamma \sim [\text{at-top}] (\lambda x. \sqrt(2*pi/x) * (x / \exp 1) \text{ powr } x :: \text{real})$ 
proof -
  note Gamma-asymp-equiv-aux
  also have eventually  $(\lambda x. \exp c * x \text{ powr } (x - 1/2) / \exp x = \sqrt(2*pi/x) * (x / \exp 1) \text{ powr } x)$  at-top
    using eventually-gt-at-top[of 0::real]
  proof eventually-elim
    fix  $x :: \text{real}$  assume  $x: x > 0$ 
    thus  $\exp c * x \text{ powr } (x - 1/2) / \exp x = \sqrt(2*pi/x) * (x / \exp 1) \text{ powr } x$ 
      by (subst powr-diff)
        (simp add: exp-c powr-half-sqrt powr-divide field-simps real-sqrt-divide)
    qed
    finally show ?thesis .
  qed

```

```

theorem fact-asymp-equiv:
   $\Gamma \sim [\text{at-top}] (\lambda n. \sqrt(2*pi*n) * (n / \exp 1) ^ n :: \text{real})$ 
proof -
  note fact-asymp-equiv-aux
  also have eventually  $(\lambda n. \exp c * \sqrt(\text{real } n) = \sqrt(2 * pi * \text{real } n))$  at-top
    using eventually-gt-at-top[of 0::nat] by eventually-elim (simp add: exp-c real-sqrt-mult)
  also have eventually  $(\lambda n. (n / \exp 1) \text{ powr } n = (n / \exp 1) ^ n)$  at-top
    using eventually-gt-at-top[of 0::nat] by eventually-elim (simp add: powr-realpow)
  finally show ?thesis .
qed

```

```

corollary stirling-tendsto-sqrt-pi:
   $(\lambda n. \text{fact } n / (\sqrt(2 * n) * (n / \exp 1) ^ n)) \longrightarrow \sqrt{\pi}$ 
proof -
  have *:  $(\lambda n. \text{fact } n / (\sqrt(2 * pi * n) * (n / \exp 1) ^ n)) \longrightarrow 1$ 
    using fact-asymp-equiv by (simp add: asymp-equiv-def)
  have  $(\lambda n. \sqrt{\pi} * \text{fact } n / (\sqrt(2 * pi * \text{real } n) * (\text{real } n / \exp 1) ^ n)) = (\lambda n. \text{fact } n / (\sqrt(\text{real } (2 * n)) * (\text{real } n / \exp 1) ^ n))$ 
    by (force simp add: divide-simps powr-realpow real-sqrt-mult)
  with tendsto-mult-left[OF *, of sqrt pi] show ?thesis by simp
qed

```

**end**

end

## 2 Complete asymptotics of the logarithmic Gamma function

**theory** *Gamma-Asymptotics*

**imports**

*HOL-Complex-Analysis.Complex-Analysis*

*Bernoulli.Bernoulli-FPS*

*Bernoulli.Periodic-Bernpoly*

*Stirling-Formula*

**begin**

### 2.1 Auxiliary Facts

**lemma** *stirling-limit-aux1*:

$((\lambda y. \ln(1 + z * \text{of-real } y) / \text{of-real } y) \longrightarrow z)$  (*at-right 0*) **for**  $z :: \text{complex}$

**proof** (*cases*  $z = 0$ )

**case** *True*

**then show** *?thesis* **by** *simp*

**next**

**case** *False*

**have**  $((\lambda y. \ln(1 + z * \text{of-real } y)) \text{ has-vector-derivative } 1 * z)$  (*at 0*)

**by** (*rule has-vector-derivative-real-field*) (*auto intro!: derivative-eq-intros*)

**then have**  $(\lambda y. (\ln(1 + z * \text{of-real } y) - \text{of-real } y * z) / \text{of-real } |y|) \xrightarrow{0} 0$

**by** (*auto simp add: has-vector-derivative-def has-derivative-def netlimit-at scaleR-conv-of-real field-simps*)

**then have**  $((\lambda y. (\ln(1 + z * \text{of-real } y) - \text{of-real } y * z) / \text{of-real } |y|) \longrightarrow 0)$

(*at-right 0*)

**by** (*rule filterlim-mono[OF - - at-le]*) *simp-all*

**also have** *?this*  $\longleftrightarrow ((\lambda y. \ln(1 + z * \text{of-real } y) / (\text{of-real } y) - z) \longrightarrow 0)$

(*at-right 0*)

**using** *eventually-at-right-less[of 0::real]*

**by** (*intro filterlim-cong refl*) (*auto elim!: eventually-mono simp: field-simps*)

**finally show** *?thesis* **by** (*simp only: LIM-zero-iff*)

**qed**

**lemma** *stirling-limit-aux2*:

$((\lambda y. y * \ln(1 + z / \text{of-real } y)) \longrightarrow z)$  *at-top* **for**  $z :: \text{complex}$

**using** *stirling-limit-aux1* [*of z*] **by** (*subst filterlim-at-top-to-right*) (*simp add: field-simps*)

**lemma** *Union-atLeastAtMost*:

**assumes**  $N > 0$

**shows**  $(\bigcup_{n \in \{0..<N\}} \{\text{real } n..\text{real } (n + 1)\}) = \{0..\text{real } N\}$

**proof** (*intro equalityI subsetI*)

**fix**  $x$  **assume**  $x: x \in \{0..\text{real } N\}$

**thus**  $x \in (\bigcup_{n \in \{0..<N\}} \{\text{real } n..\text{real } (n + 1)\})$

**proof** (*cases*  $x = \text{real } N$ )

```

case True
with assms show ?thesis by (auto intro!: bexI[of - N - 1])
next
case False
with x have x: x ≥ 0 x < real N by simp-all
hence x ≥ real (nat ⌊x⌋) x ≤ real (nat ⌊x⌋ + 1) by linarith+
moreover from x have nat ⌊x⌋ < N by linarith
ultimately have ∃n∈{0..<N}. x ∈ {real n..real (n + 1)}
by (intro bexI[of - nat ⌊x⌋]) simp-all
thus ?thesis by blast
qed
qed auto

```

## 2.2 Cones in the complex plane

```

definition complex-cone :: real ⇒ real ⇒ complex set where
  complex-cone a b = {z. ∃y∈{a..b}. z = rcis (norm z) y}

```

```

abbreviation complex-cone' :: real ⇒ complex set where
  complex-cone' a ≡ complex-cone (-a) a

```

```

lemma zero-in-complex-cone [simp, intro]: a ≤ b ⇒ 0 ∈ complex-cone a b
by (auto simp: complex-cone-def)

```

```

lemma complex-coneE:
  assumes z ∈ complex-cone a b
  obtains r α where r ≥ 0 α ∈ {a..b} z = rcis r α
proof –
  from assms obtain y where y ∈ {a..b} z = rcis (norm z) y
    unfolding complex-cone-def by auto
  thus ?thesis using that[of norm z y] by auto
qed

```

```

lemma arg-cis [simp]:
  assumes x ∈ {-pi<..pi}
  shows Arg (cis x) = x
  using assms by (intro cis-Arg-unique) auto

```

```

lemma arg-mult-of-real-left [simp]:
  assumes r > 0
  shows Arg (of-real r * z) = Arg z
proof (cases z = 0)
  case False
  thus ?thesis
    using Arg-bounded[of z] assms
    by (intro cis-Arg-unique) (auto simp: sgn-mult sgn-of-real cis-Arg)
qed auto

```

```

lemma arg-mult-of-real-right [simp]:

```

```

assumes  $r > 0$ 
shows  $\operatorname{Arg}(z * \text{of-real } r) = \operatorname{Arg} z$ 
by (subst mult.commute, subst arg-mult-of-real-left) (simp-all add: assms)

lemma arg-rcis [simp]:
assumes  $x \in \{-pi <.. pi\}$   $r > 0$ 
shows  $\operatorname{Arg}(\operatorname{rcis} r x) = x$ 
using assms by (simp add: rcis-def)

lemma rcis-in-complex-cone [intro]:
assumes  $\alpha \in \{a..b\}$   $r \geq 0$ 
shows  $\operatorname{rcis} r \alpha \in \text{complex-cone } a b$ 
using assms by (auto simp: complex-cone-def)

lemma arg-imp-in-complex-cone:
assumes  $\operatorname{Arg} z \in \{a..b\}$ 
shows  $z \in \text{complex-cone } a b$ 
proof -
have  $z = \operatorname{rcis}(\operatorname{norm} z) (\operatorname{Arg} z)$ 
by (simp add: rcis-cmod-Arg)
also have ...  $\in \text{complex-cone } a b$ 
using assms by auto
finally show ?thesis .
qed

lemma complex-cone-altdef:
assumes  $-pi < a$   $a \leq b$   $b \leq pi$ 
shows  $\text{complex-cone } a b = \operatorname{insert} 0 \{z. \operatorname{Arg} z \in \{a..b\}\}$ 
proof (intro equalityI subsetI)
fix  $z$  assume  $z \in \text{complex-cone } a b$ 
then obtain  $r \alpha$  where  $*: r \geq 0$   $\alpha \in \{a..b\}$   $z = \operatorname{rcis} r \alpha$ 
by (auto elim: complex-coneE)
have  $\operatorname{Arg} z \in \{a..b\}$  if [simp]:  $z \neq 0$ 
proof -
have  $r > 0$  using that  $*$  by (subst (asm) *) auto
hence  $\alpha \in \{a..b\}$ 
using *(1,2) assms by (auto simp: *(1))
moreover from assms *(2) have  $\alpha \in \{-pi <.. pi\}$ 
by auto
ultimately show ?thesis using *(3) { $r > 0$ }
by (subst *) auto
qed
thus  $z \in \operatorname{insert} 0 \{z. \operatorname{Arg} z \in \{a..b\}\}$ 
by auto
qed (use assms in {auto intro: arg-imp-in-complex-cone})

lemma nonneg-of-real-in-complex-cone [simp, intro]:
assumes  $x \geq 0$   $a \leq 0$   $0 \leq b$ 
shows  $\text{of-real } x \in \text{complex-cone } a b$ 

```

```

proof -
  from assms have rcis x 0 ∈ complex-cone a b
    by (intro rcis-in-complex-cone) auto
    thus ?thesis by simp
  qed

lemma one-in-complex-cone [simp, intro]: a ≤ 0 ⇒ 0 ≤ b ⇒ 1 ∈ complex-cone
a b
  using nonneg-of-real-in-complex-cone[of 1] by (simp del: nonneg-of-real-in-complex-cone)

lemma of-nat-in-complex-cone [simp, intro]: a ≤ 0 ⇒ 0 ≤ b ⇒ of-nat n ∈
complex-cone a b
  using nonneg-of-real-in-complex-cone[of real n] by (simp del: nonneg-of-real-in-complex-cone)

```

## 2.3 Another integral representation of the Beta function

```

lemma complex-cone-inter-nonpos-Reals:
  assumes -pi < a a ≤ b b < pi
  shows complex-cone a b ∩ ℝ≤0 = {0}
proof (safe elim!: nonpos-Reals-cases)
  fix x :: real
  assume complex-of-real x ∈ complex-cone a b x ≤ 0
  hence ¬(x < 0)
    using assms by (intro notI) (auto simp: complex-cone-altdef)
    with ⟨x ≤ 0⟩ show complex-of-real x = 0 by auto
  qed (use assms in auto)

```

```

theorem
  assumes a: a > 0 and b: b > (0 :: real)
  shows has-integral-Beta-real':
    ((λu. u powr (b - 1) / (1 + u) powr (a + b)) has-integral Beta a b) {0 < ..}
  and Beta-conv-nn-integral:
    Beta a b = (ʃ+u. ennreal (indicator {0 < ..} u * u powr (b - 1) / (1 +
u) powr (a + b)) ∂lborel)
proof -
  define I where
    I = (ʃ+u. ennreal (indicator {0 < ..} u * u powr (b - 1) / (1 + u) powr (a +
b)) ∂lborel)
  have Gamma (a + b) > 0 Beta a b > 0
    using assms by (simp-all add: add-pos-pos Beta-def)
  from a b have ennreal (Gamma a * Gamma b) =
    (ʃ+t. ennreal (indicator {0..} t * t powr (a - 1) / exp t) ∂lborel) *
    (ʃ+t. ennreal (indicator {0..} t * t powr (b - 1) / exp t) ∂lborel)
    by (subst ennreal-mult') (simp-all add: Gamma-conv-nn-integral-real)
  also have ... = (ʃ+t. ʃ+u. ennreal (indicator {0..} t * t powr (a - 1) / exp
t) *
    ennreal (indicator {0..} u * u powr (b - 1) / exp u) ∂lborel
  ∂lborel
    by (simp add: nn-integral-cmult nn-integral-multc)

```

```

also have ... = ( $\int^+ t. \text{indicator } \{0 <..\} t * (\int^+ u. \text{indicator } \{0..\} u * t \text{powr } (a - 1) * u \text{powr } (b - 1)$ 
 $\quad / \exp(t + u) \partial\text{borel}) \partial\text{borel}$ )
by (intro nn-integral-cong-AE AE-I[of - - {0}])
(auto simp: indicator-def divide-ennreal ennreal-mult' [symmetric] exp-add mult-ac)
also have ... = ( $\int^+ t. \text{indicator } \{0 <..\} t * (\int^+ u. \text{indicator } \{0..\} u * t \text{powr } (a - 1) * u \text{powr } (b - 1)$ 
 $\quad / \exp(t + u)$ 
 $\quad \partial(\text{density } (\text{distr lborel borel } ((*) t)) (\lambda x. \text{ennreal } |t|))) \partial\text{borel}$ )
by (intro nn-integral-cong mult-indicator-cong, subst lborel-distr-mult' [symmetric])
auto
also have ... = ( $\int^+(t:\text{real}). \text{indicator } \{0 <..\} t * (\int^+ u.$ 
 $\quad \text{indicator } \{0..\} (u * t) * t \text{powr } a *$ 
 $\quad (u * t) \text{powr } (b - 1) / \exp(t + t * u) \partial\text{borel}) \partial\text{borel}$ )
by (intro nn-integral-cong mult-indicator-cong)
(auto simp: nn-integral-density nn-integral-distr algebra-simps powr-diff
simp flip: ennreal-mult)
also have ... = ( $\int^+(t:\text{real}). \int^+ u. \text{indicator } (\{0 <..\} \times \{0..\}) (t, u) *$ 
 $\quad t \text{powr } a * (u * t) \text{powr } (b - 1) / \exp(t * (u + 1)) \partial\text{borel} \partial\text{borel}$ )
by (subst nn-integral-cmult [symmetric], simp, intro nn-integral-cong)
(auto simp: indicator-def zero-le-mult-iff algebra-simps)
also have ... = ( $\int^+(t:\text{real}). \int^+ u. \text{indicator } (\{0 <..\} \times \{0..\}) (t, u) *$ 
 $\quad t \text{powr } (a + b - 1) * u \text{powr } (b - 1) / \exp(t * (u + 1)) \partial\text{borel}$ 
 $\partial\text{borel}$ )
by (intro nn-integral-cong) (auto simp: powr-add powr-diff indicator-def powr-mult field-simps)
also have ... = ( $\int^+(u:\text{real}). \int^+ t. \text{indicator } (\{0 <..\} \times \{0..\}) (t, u) *$ 
 $\quad t \text{powr } (a + b - 1) * u \text{powr } (b - 1) / \exp(t * (u + 1)) \partial\text{borel}$ 
 $\partial\text{borel}$ )
by (rule lborel-pair.Fubini') auto
also have ... = ( $\int^+(u:\text{real}). \text{indicator } \{0..\} u * (\int^+ t. \text{indicator } \{0 <..\} t *$ 
 $\quad t \text{powr } (a + b - 1) * u \text{powr } (b - 1) / \exp(t * (u + 1)) \partial\text{borel}$ 
 $\partial\text{borel}$ )
by (intro nn-integral-cong mult-indicator-cong) (auto simp: indicator-def)
also have ... = ( $\int^+(u:\text{real}). \text{indicator } \{0 <..\} u * (\int^+ t. \text{indicator } \{0 <..\} t *$ 
 $\quad t \text{powr } (a + b - 1) * u \text{powr } (b - 1) / \exp(t * (u + 1)) \partial\text{borel}$ 
 $\partial\text{borel}$ )
by (intro nn-integral-cong-AE AE-I[of - - {0}]) (auto simp: indicator-def)
also have ... = ( $\int^+(u:\text{real}). \text{indicator } \{0 <..\} u * (\int^+ t. \text{indicator } \{0 <..\} t *$ 
 $\quad t \text{powr } (a + b - 1) * u \text{powr } (b - 1) / \exp(t * (u + 1))$ 
 $\quad \partial(\text{density } (\text{distr lborel borel } ((*) (1/(1+u)))) (\lambda x. \text{ennreal } |1/(1+u)|))) \partial\text{borel}$ )
by (intro nn-integral-cong mult-indicator-cong, subst lborel-distr-mult' [symmetric])
auto
also have ... = ( $\int^+(u:\text{real}). \text{indicator } \{0 <..\} u *$ 
 $\quad (\int^+ t. \text{ennreal } (1 / (u + 1)) * \text{ennreal } (\text{indicator } \{0 <..\} (t / (u + 1)) *$ 
 $\quad (t / (1+u)) \text{powr } (a + b - 1) * u \text{powr } (b - 1) / \exp t)$ 

```

```

 $\partial lborel) \partial lborel)$ 
by (intro nn-integral-cong mult-indicator-cong)
  (auto simp: nn-integral-distr nn-integral-density add-ac)
also have ... = ( $\int^+ u. \int^+ t. indicator (\{0 <..\} \times \{0 <..\}) (u, t) * \frac{1}{(u+1)} * (t / (u+1)) powr (a + b - 1) * u powr (b - 1) / exp t$ )
   $\partial lborel \partial lborel)$ 
by (subst nn-integral-cmult [symmetric], simp, intro nn-integral-cong)
  (auto simp: indicator-def field-simps divide-ennreal simp flip: ennreal-mult ennreal-mult')
also have ... = ( $\int^+ u. \int^+ t. ennreal (indicator \{0 <..\} u * u powr (b - 1) / (1 + u) powr (a + b)) * ennreal (indicator \{0 <..\} t * t powr (a + b - 1) / exp t$ )
   $\partial lborel \partial lborel)$ 
by (intro nn-integral-cong)
  (auto simp: indicator-def powr-add powr-diff powr-divide powr-minus divide-simps add-ac
    simp flip: ennreal-mult)
also have ... =  $I * (\int^+ t. indicator \{0 <..\} t * t powr (a + b - 1) / exp t$ )
 $\partial lborel)$ 
  by (simp add: nn-integral-cmult nn-integral-multc I-def)
also have ( $\int^+ t. indicator \{0 <..\} t * t powr (a + b - 1) / exp t \partial lborel$ ) =
  ennreal (Gamma (a + b))
using assms
by (subst Gamma-conv-nn-integral-real)
  (auto intro!: nn-integral-cong-AE[OF AE-I[of - - {0}]] simp: indicator-def split: if-splits split-of-bool-asm)
finally have ennreal (Gamma a * Gamma b) =  $I * ennreal (\Gamma (a + b))$  .
hence ennreal (Gamma a * Gamma b) / ennreal (Gamma (a + b)) =
   $I * ennreal (\Gamma (a + b)) / ennreal (\Gamma (a + b))$  by simp
also have ... =  $I$ 
  using <math>\Gamma (a + b) > 0</math> by (intro ennreal-mult-divide-eq) auto
also have ennreal (Gamma a * Gamma b) / ennreal (Gamma (a + b)) =
  ennreal (Gamma a * Gamma b / Gamma (a + b))
  using assms by (intro divide-ennreal) auto
also have ... = ennreal (Beta a b)
  by (simp add: Beta-def)
finally show *: ennreal (Beta a b) =  $I$  .

define f where f =  $(\lambda u. u powr (b - 1) / (1 + u) powr (a + b))$ 
have (( $\lambda u. indicator \{0 <..\} u * f u$ ) has-integral Beta a b) UNIV
  using * <math>\Gamma a b > 0</math>
  by (subst has-integral-iff-nn-integral-lebesgue)
    (auto simp: f-def measurable-completion nn-integral-completion I-def mult-ac)
also have ( $\lambda u. indicator \{0 <..\} u * f u$ ) =  $(\lambda u. if u \in \{0 <..\} then f u else 0)$ 
  by (auto simp: fun-eq-iff)
also have (... has-integral Beta a b) UNIV  $\longleftrightarrow$  (f has-integral Beta a b) {0 <..}
  by (rule has-integral-restrict-UNIV)
finally show ... by (simp add: f-def)
qed
```

```

lemma has-integral-Beta2:
  fixes a :: real
  assumes a < -1/2
  shows ((λx. (1 + x ^ 2) powr a) has-integral Beta (- a - 1 / 2) (1 / 2) / 2) {0<..}
proof -
  define f where f = (λu. u powr (-1/2) / (1 + u) powr (-a))
  define C where C = Beta (-a-1/2) (1/2)
  have I: (f has-integral C) {0<..}
    using has-integral-Beta-real'[of -a-1/2 1/2] assms
    by (simp-all add: diff-divide-distrib f-def C-def)

  define g where g = (λx. x ^ 2 :: real)
  have bij: bij-betw g {0<..} {0<..}
    by (intro bij-betwI[of --- sqrt]) (auto simp: g-def)

  have (f absolutely-integrable-on g ` {0<..} ∧ integral (g ` {0<..}) f = C)
    using I bij by (simp add: bij-betw-def has-integral-iff absolutely-integrable-on-def f-def)
  also have ?this ←→ ((λx. |2 * x| *R f (g x)) absolutely-integrable-on {0<..} ∧
    integral {0<..} (λx. |2 * x| *R f (g x)) = C)
    using bij by (intro has-absolute-integral-change-of-variables-1' [symmetric])
    (auto intro!: derivative-eq-intros simp: g-def bij-betw-def)
  finally have ((λx. |2 * x| * f (g x)) has-integral C) {0<..}
    by (simp add: absolutely-integrable-on-def f-def has-integral-iff)
  also have ?this ←→ ((λx::real. 2 * (1 + x^2) powr a) has-integral C) {0<..}
    by (intro has-integral-cong) (auto simp: f-def g-def powr-def exp-minus ln-realpow field-simps)
  finally have ((λx::real. 1/2 * (2 * (1 + x^2) powr a)) has-integral 1/2 * C)
    {0<..}
    by (intro has-integral-mult-right)
  thus ?thesis by (simp add: C-def)
qed

lemma has-integral-Beta3:
  fixes a b :: real
  assumes a < -1/2 and b > 0
  shows ((λx. (b + x ^ 2) powr a) has-integral
    Beta (-a - 1/2) (1/2) / 2 * b powr (a + 1/2)) {0<..}
proof -
  define C where C = Beta (- a - 1 / 2) (1 / 2) / 2
  have int: nn-integral_lborel (λx. indicator {0<..} x * (1 + x ^ 2) powr a) = C
    using nn-integral-has-integral-lebesgue[OF - has-integral-Beta2[OF assms(1)]]
    by (auto simp: C-def)
  have nn-integral_lborel (λx. indicator {0<..} x * (b + x ^ 2) powr a) =
    (ʃ+ x. ennreal (indicat-real {0<..} (x * sqrt b) * (b + (x * sqrt b)^2) powr a
    * sqrt b) ∂lborel)
    using assms

```

```

by (subst lborel-distr-mult[of sqrt b])
  (auto simp: nn-integral-density nn-integral-distr mult-ac simp flip: ennreal-mult)
also have ... = ( $\int^+ x.$  ennreal (indicat-real {0<..} x * (b * (1 + x ^ 2)) powr a * sqrt b)  $\partial$ borel)
  using assms
by (intro nn-integral-cong) (auto simp: indicator-def field-simps zero-less-mult-iff)
also have ... = ( $\int^+ x.$  ennreal (indicat-real {0<..} x * b powr (a + 1/2) * (1 + x ^ 2) powr a)  $\partial$ borel)
  using assms
by (intro nn-integral-cong) (auto simp: indicator-def powr-add powr-half-sqrt powr-mult)
also have ... = b powr (a + 1/2) * ( $\int^+ x.$  ennreal (indicat-real {0<..} x * (1 + x ^ 2) powr a)  $\partial$ borel)
  using assms by (subst nn-integral-cmult [symmetric]) (simp-all add: mult-ac flip: ennreal-mult)
also have ( $\int^+ x.$  ennreal (indicat-real {0<..} x * (1 + x ^ 2) powr a)  $\partial$ borel)
= C
  using int by simp
also have ennreal (b powr (a + 1/2)) * ennreal C = ennreal (C * b powr (a + 1/2))
  using assms by (subst ennreal-mult) (auto simp: C-def mult-ac Beta-def)
finally have *: ( $\int^+ x.$  ennreal (indicat-real {0<..} x * (b + x ^ 2) powr a)  $\partial$ borel)
= ...
  hence (( $\lambda x.$  indicator {0<..} x * (b + x ^ 2) powr a) has-integral C * b powr (a + 1/2)) UNIV
  using assms
by (subst has-integral-iff-nn-integral-lebesgue)
  (auto simp: C-def measurable-completion nn-integral-completion Beta-def)
also have ( $\lambda x.$  indicator {0<..} x * (b + x ^ 2) powr a) =
  ( $\lambda x.$  if x ∈ {0<..} then (b + x ^ 2) powr a else 0)
  by (auto simp: fun-eq-iff)
finally show ?thesis
  by (subst (asm) has-integral-restrict-UNIV) (auto simp: C-def)
qed

```

## 2.4 Asymptotics of the real $\log \Gamma$ function and its derivatives

This is the error term that occurs in the expansion of *ln-Gamma*. It can be shown to be of order  $O(s^{-n})$ .

**definition** stirling-integral :: nat ⇒ 'a :: {real-normed-div-algebra, banach} ⇒ 'a  
**where**

stirling-integral n s =  
 $\lim (\lambda N. \text{integral } \{0..N\} (\lambda x. \text{of-real } (\text{pbernpoly } n x) / (\text{of-real } x + s) ^ n))$

**context**

fixes s :: complex **assumes** s:  $s \notin \mathbb{R}_{\leq 0}$

fixes approx :: nat ⇒ complex

defines approx ≡ ( $\lambda N.$

$(\sum n = 1..N. s / \text{of-nat } n - \ln(1 + s / \text{of-nat } n)) - (\text{euler-mascheroni} * s)$

```

+ ln s) — — → ln-Gamma s
  (ln-Gamma (of-nat N) - ln (2 * pi / of-nat N) / 2 - of-nat N * ln (of-nat
N) + of-nat N) — — → 0
  s * (harm (N - 1) - ln (of-nat (N - 1)) - euler-mascheroni) + — → 0
  s * (ln (of-nat N + s) - ln (of-nat (N - 1))) — — → 0
  (1/2) * (ln (of-nat N + s) - ln (of-nat N)) + — — → 0
  of-nat N * (ln (of-nat N + s) - ln (of-nat N)) — — → s
  (s - 1/2) * ln s - ln (2 * pi) / 2

begin

qualified lemma
assumes N: N > 0
shows integrable-pbernpoly-1:
  ( $\lambda x. \text{of-real}(-\text{pbernpoly } 1 x) / (\text{of-real } x + s)$ ) integrable-on {0..real N}
and integral-pbernpoly-1-aux:
  integral {0..real N} ( $\lambda x. -\text{of-real}(\text{pbernpoly } 1 x) / (\text{of-real } x + s)$ ) =
approx N
and has-integral-pbernpoly-1:
  (( $\lambda x. \text{pbernpoly } 1 x / (x + s)$ ) has-integral
  ( $\sum m < N. (\text{of-nat } m + 1 / 2 + s) * (\ln(\text{of-nat } m + s) - \ln(\text{of-nat } m + 1 + s)) + 1$ )) {0..real N}

proof -
let ?A = ( $\lambda n. \{\text{of-nat } n .. \text{of-nat } (n + 1)\}$ ) ` {0..< N}
have has-integral:
  (( $\lambda x. -\text{pbernpoly } 1 x / (x + s)$ ) has-integral
  ( $(\text{of-nat } n + 1/2 + s) * (\ln(\text{of-nat } (n + 1) + s) - \ln(\text{of-nat } n + s)) - 1$ )
  {of-nat n..of-nat (n + 1)} for n
proof (rule has-integral-spike)
have (( $\lambda x. (\text{of-nat } n + 1/2 + s) * (1 / (\text{of-real } x + s)) - 1$ ) has-integral
  ( $(\text{of-nat } n + 1/2 + s) * (\ln(\text{of-real } (\text{real } (n + 1)) + s) - \ln(\text{of-real } (\text{real } n) + s)) - 1$ )
  {of-nat n..of-nat (n + 1)})
using s has-integral-const-real[of 1 of-nat n of-nat (n + 1)]
by (intro has-integral-diff has-integral-mult-right fundamental-theorem-of-calculus)
(auto intro!: derivative-eq-intros has-vector-derivative-real-field
simp: has-real-derivative-iff-has-vector-derivative [symmetric] field-simps
complex-nonpos-Reals-iff)
thus (( $\lambda x. (\text{of-nat } n + 1/2 + s) * (1 / (\text{of-real } x + s)) - 1$ ) has-integral
  ( $(\text{of-nat } n + 1/2 + s) * (\ln(\text{of-nat } (n + 1) + s) - \ln(\text{of-nat } n + s)) - 1$ )
  {of-nat n..of-nat (n + 1)}) by simp

show -pbernpoly 1 x / (x + s) = ( $\text{of-nat } n + 1/2 + s$ ) * (1 / (x + s)) - 1
if x ∈ {of-nat n..of-nat (n + 1)} - {of-nat (n + 1)} for x
proof -
have x: x ≥ real n x < real (n + 1) using that by simp-all
hence floor x = int n by linarith
moreover from s x have complex-of-real x ≠ -s

```

```

by (auto simp add: complex-eq-iff complex-nonpos-Reals-iff simp del: of-nat-Suc)
ultimately show -pbernpoly 1 x / (x + s) = (of-nat n + 1/2 + s) * (1 /
(x + s)) - 1
  by (auto simp: pbernpoly-def bernpoly-def frac-def divide-simps add-eq-0-iff2)
qed
qed simp-all
hence *:  $\bigwedge I. I \in ?A \implies ((\lambda x. -pbernpoly 1 x / (x + s)) \text{ has-integral}$ 
 $(\ln(\inf I + 1/2 + s) * (\ln(\inf I + 1 + s) - \ln(\inf I + s)) - 1) I$ 
  by (auto simp: add-ac)
have  $((\lambda x. -pbernpoly 1 x / (x + s)) \text{ has-integral}$ 
 $(\sum I \in ?A. (\inf I + 1/2 + s) * (\ln(\inf I + 1 + s) - \ln(\inf I + s)) -$ 
 $1))$ 
   $(\bigcup_{n \in \{0..<N\}} \{\text{real } n.. \text{real } (n + 1)\})$  (is (- has-integral ?i) -)
apply (intro has-integral-Union * finite-imageI)
apply (force intro!: negligible-atLeastAtMostI pairwiseI) +
done
hence has-integral:  $((\lambda x. -pbernpoly 1 x / (x + s)) \text{ has-integral ?i}) \{0.. \text{real } N\}$ 
  by (subst has-integral-spoke-set-eq)
  (use Union-atLeastAtMost assms in ⟨auto simp: intro!: empty-imp-negligible⟩)
hence  $(\lambda x. -pbernpoly 1 x / (x + s)) \text{ integrable-on } \{0.. \text{real } N\}$ 
  and integral:  $\text{integral } \{0.. \text{real } N\} (\lambda x. -pbernpoly 1 x / (x + s)) = ?i$ 
  by (simp-all add: has-integral-iff)
show  $(\lambda x. -pbernpoly 1 x / (x + s)) \text{ integrable-on } \{0.. \text{real } N\}$  by fact

note has-integral-neg[OF has-integral]
also have  $-?i = (\sum x < N. (\text{of-nat } x + 1/2 + s) * (\ln(\text{of-nat } x + s) - \ln(\text{of-nat } x + 1 + s)) + 1)$ 
  by (subst sum.reindex)
  (simp-all add: inj-on-def atLeast0LessThan algebra-simps sum-negf [symmetric])
finally show has-integral:
   $((\lambda x. \text{of-real } (pbernpoly 1 x) / (\text{of-real } x + s)) \text{ has-integral } \dots) \{0.. \text{real } N\}$  by
simp

note integral
also have ?i =  $(\sum n < N. (\text{of-nat } n + 1/2 + s) *$ 
 $(\ln(\text{of-nat } n + 1 + s) - \ln(\text{of-nat } n + s))) - N$  (is - = ?S - -)
  by (subst sum.reindex) (simp-all add: inj-on-def sum-subtractf atLeast0LessThan)
also have ?S =  $(\sum n < N. \text{of-nat } n * (\ln(\text{of-nat } n + 1 + s) - \ln(\text{of-nat } n +$ 
 $s))) +$ 
 $(s + 1/2) * (\sum n < N. \ln(\text{of-nat } (\text{Suc } n) + s) - \ln(\text{of-nat } n +$ 
 $s))$ 
  (is - = ?S1 + - * ?S2) by (simp add: algebra-simps sum.distrib sum-subtractf
sum-distrib-left)
also have ?S2 =  $\ln(\text{of-nat } N + s) - \ln s$  by (subst sum-lessThan-telescope)
simp
also have ?S1 =  $(\sum n=1..< N. \text{of-nat } n * (\ln(\text{of-nat } n + 1 + s) - \ln(\text{of-nat } n + s)))$ 
  by (intro sum.mono-neutral-right) auto
also have ... =  $(\sum n=1..< N. \text{of-nat } n * \ln(\text{of-nat } n + 1 + s)) - (\sum n=1..< N.$ 

```

```

 $of\text{-}nat n * ln (of\text{-}nat n + s))$ 
  by (simp add: algebra-simps sum-subtractf)
also have  $(\sum_{n=1..<N} of\text{-}nat n * ln (of\text{-}nat n + 1 + s)) =$ 
   $(\sum_{n=1..<N} (of\text{-}nat n - 1) * ln (of\text{-}nat n + s)) + (N - 1) * ln$ 
 $(of\text{-}nat N + s)$ 
  by (induction N) (simp-all add: add-ac of-nat-diff)
also have ...  $- (\sum_{n=1..<N} of\text{-}nat n * ln (of\text{-}nat n + s)) =$ 
   $- (\sum_{n=1..<N} ln (of\text{-}nat n + s)) + (N - 1) * ln (of\text{-}nat N + s)$ 
  by (induction N) (simp-all add: algebra-simps)
also from s have neq:  $s + of\text{-}nat x \neq 0$  for x
  by (auto simp: complex-nonpos-Reals-iff complex-eq-iff)
hence  $(\sum_{n=1..<N} ln (of\text{-}nat n + s)) = (\sum_{n=1..<N} ln (of\text{-}nat n) + ln (1$ 
 $+ s/n))$ 
  by (intro sum.cong refl, subst Ln-times-of-nat [symmetric]) (auto simp: di-
vide-simps add-ac)
also have ...  $= ln (fact (N - 1)) + (\sum_{n=1..<N} ln (1 + s/n))$ 
  by (induction N) (simp-all add: Ln-times-of-nat fact-reduce add-ac)
also have  $(\sum_{n=1..<N} ln (1 + s/n)) = -(\sum_{n=1..<N} s / n - ln (1 + s/n))$ 
 $+ s * (\sum_{n=1..<N} 1 / of\text{-}nat n)$ 
  by (simp add: sum-distrib-left sum-subtractf)
also from N have  $ln (fact (N - 1)) = ln\text{-}\Gamma (of\text{-}nat N :: complex)$ 
  by (simp add: ln-Gamma-complex-conv-fact)
also have  $\{1..<N\} = \{1..N - 1\}$  by auto
hence  $(\sum_{n=1..<N} 1 / of\text{-}nat n) = (harm (N - 1) :: complex)$ 
  by (simp add: harm-def divide-simps)
also have  $- (ln\text{-}\Gamma (of\text{-}nat N)) + (- (\sum_{n=1..<N} s / of\text{-}nat n - ln (1$ 
 $+ s / of\text{-}nat n)) +$ 
   $s * harm (N - 1))) + of\text{-}nat (N - 1) * ln (of\text{-}nat N + s) +$ 
   $(s + 1 / 2) * (ln (of\text{-}nat N + s) - ln s) - of\text{-}nat N = approx N$ 
using N by (simp add: field-simps of-nat-diff ln-div approx-def Ln-of-nat
  ln-Gamma-complex-of-real [symmetric])
finally show integral  $\{0..of\text{-}nat N\} (\lambda x. -of\text{-}real (pbernpoly 1 x) / (of\text{-}real x +$ 
 $s)) = \dots$ 
  by simp
qed

```

```

lemma integrable-ln-Gamma-aux:
shows  $(\lambda x. of\text{-}real (pbernpoly n x) / (of\text{-}real x + s) ^ n)$  integrable-on  $\{0..real N\}$ 
proof (cases n = 1)
  case True
  with s show ?thesis using integrable-neg[OF integrable-pbernpoly-1[of N]]
    by (cases N = 0) (simp-all add: integrable-negligible)
next
  case False
  from s have  $of\text{-}real x + s \neq 0$  if  $x \geq 0$  for x using that
    by (auto simp: complex-eq-iff add-eq-0-iff2 complex-nonpos-Reals-iff)
  with False s show ?thesis
    by (auto intro!: integrable-continuous-real continuous-intros)

```

**qed**

This following proof is based on “Rudiments of the theory of the gamma function” by Bruce Berndt [1].

```

lemma tendsto-of-real-0-I:
  ( $f \rightarrow 0$ )  $G \implies ((\lambda x. (\text{of-real } (f x))) \rightarrow (0 ::'a::\text{real-normed-div-algebra}))$ 
   $G$ 
  using tendsto-of-real-iff by force

qualified lemma integral-pbernpoly-1:
   $(\lambda N. \text{integral } \{0..\text{real } N\} (\lambda x. \text{pbernpoly } 1 x / (x + s))) \rightarrow -\ln\text{-Gamma } s - s + (s - 1 / 2) * \ln s + \ln(2 * \pi) / 2$ 
proof -
  have neq:  $s + \text{of-real } x \neq 0$  if  $x \geq 0$  for  $x :: \text{real}$ 
    using that  $s$  by (auto simp: complex-eq-iff complex-nonpos-Reals-iff)
  have (approx  $\rightarrow \ln\text{-Gamma } s - 0 - 0 + 0 - 0 + s - (s - 1/2) * \ln s - \ln(2 * \pi) / 2$ ) at-top
    unfolding approx-def
  proof (intro tendsto-add tendsto-diff)
    from  $s$  have  $s' : s \notin \mathbb{Z}_{\leq 0}$  by (auto simp: complex-nonpos-Reals-iff elim!: non-pos-Ints-cases)
    have  $(\lambda n. \sum_{i=1..n} s / \text{of-nat } i - \ln(1 + s / \text{of-nat } i)) \rightarrow \ln\text{-Gamma } s + \text{euler-mascheroni} * s + \ln s$  (is ?f  $\rightarrow \dots$ )
      using ln-Gamma-series'-aux[ $\text{OF } s'$ ] unfolding sums-def
      by (subst filterlim-sequentially-Suc [symmetric], subst (asm) sum.atLeast1-atMost-eq [symmetric])
        (simp add: atLeastLessThanSuc-atLeastAtMost)
    thus  $((\lambda n. ?f n - (\text{euler-mascheroni} * s + \ln s)) \rightarrow \ln\text{-Gamma } s)$  at-top
      by (auto intro: tendsto-eq-intros)
  next
    show  $(\lambda x. \text{complex-of-real } (\ln\text{-Gamma } (\text{real } x) - \ln(2 * \pi / \text{real } x) / 2 - \text{real } x * \ln(\text{real } x) + \text{real } x)) \rightarrow 0$ 
  proof (intro tendsto-of-real-0-I)
    filterlim-compose[ $\text{OF tendsto-sandwich filterlim-real-sequentially}]$ 
    show eventually  $(\lambda x :: \text{real}. \ln\text{-Gamma } x - \ln(2 * \pi / x) / 2 - x * \ln x + x \geq 0)$  at-top
      using eventually-ge-at-top[of 1::real]
      by eventually-elim (insert ln-Gamma-bounds(1), simp add: algebra-simps)
    show eventually  $(\lambda x :: \text{real}. \ln\text{-Gamma } x - \ln(2 * \pi / x) / 2 - x * \ln x + x \leq 1 / 12 * \text{inverse } x)$  at-top
      using eventually-ge-at-top[of 1::real]
      by eventually-elim (insert ln-Gamma-bounds(2), simp add: field-simps)
    show  $((\lambda x :: \text{real}. 1 / 12 * \text{inverse } x) \rightarrow 0)$  at-top
      by real-asymp
    qed simp-all
  next
    have  $(\lambda x. s * \text{of-real } (\text{harm } (x - 1) - \ln(\text{real } (x - 1)) - \text{euler-mascheroni})) \rightarrow$ 

```

```

 $s * \text{of-real}(\text{euler-mascheroni} - \text{euler-mascheroni})$ 
by (subst filterlim-sequentially-Suc [symmetric], intro tendsto-intros)
      (insert euler-mascheroni-LIMSEQ, simp-all)
also have ?this  $\longleftrightarrow (\lambda x. s * (\text{harm}(x - 1) - \ln(\text{of-nat}(x - 1)) - \text{euler-mascheroni})) \longrightarrow 0$ 
by (intro filterlim-cong refl eventually-mono[OF eventually-gt-at-top[of 1::nat]])
      (auto simp: of-real-harm simp del: of-nat-diff)
finally show  $(\lambda x. s * (\text{harm}(x - 1) - \ln(\text{of-nat}(x - 1)) - \text{euler-mascheroni})) \longrightarrow 0$  .

next
have  $((\lambda x. \ln(1 + (s + 1) / \text{of-real } x)) \longrightarrow \ln(1 + 0)) \text{ at-top (is ?P)}$ 
by (intro tendsto-intros tendsto-divide-0[OF tendsto-const])
      (simp-all add: filterlim-ident filterlim-at-infinity-conv-norm-at-top filterlim-abs-real)
also have  $\ln(\text{of-real}(x + 1) + s) - \ln(\text{complex-of-real } x) = \ln(1 + (s + 1) / \text{of-real } x)$ 
if  $x > 1$  for  $x$  using that  $s$ 
using Ln-divide-of-real[of  $x$  of-real  $(x + 1) + s$ , symmetric] neq[of  $x + 1$ ]
by (simp add: field-simps Ln-of-real)
hence ?P  $\longleftrightarrow ((\lambda x. \ln(\text{of-real}(x + 1) + s) - \ln(\text{of-real } x)) \longrightarrow 0) \text{ at-top}$ 
by (intro filterlim-cong refl)
      (auto intro: eventually-mono[OF eventually-gt-at-top[of 1::real]])
finally have  $((\lambda n. \ln(\text{of-real}(real n + 1) + s) - \ln(\text{of-real}(real n))) \longrightarrow 0) \text{ at-top}$ 
by (rule filterlim-compose[OF - filterlim-real-sequentially])
hence  $((\lambda n. \ln(\text{of-nat } n + s) - \ln(\text{of-nat}(n - 1))) \longrightarrow 0) \text{ at-top}$ 
by (subst filterlim-sequentially-Suc [symmetric]) (simp add: add-ac)
thus  $(\lambda x. s * (\ln(\text{of-nat } x + s) - \ln(\text{of-nat}(x - 1)))) \longrightarrow 0$ 
by (rule tendsto-mult-right-zero)

next
have  $((\lambda x. \ln(1 + s / \text{of-real } x)) \longrightarrow \ln(1 + 0)) \text{ at-top (is ?P)}$ 
by (intro tendsto-intros tendsto-divide-0[OF tendsto-const])
      (simp-all add: filterlim-ident filterlim-at-infinity-conv-norm-at-top filterlim-abs-real)
also have  $\ln(\text{of-real } x + s) - \ln(\text{of-real } x) = \ln(1 + s / \text{of-real } x)$  if  $x > 0$ 
for  $x$ 
using Ln-divide-of-real[of  $x$  of-real  $x + s$ ] neq[of  $x$ ] that
by (auto simp: field-simps Ln-of-real)
hence ?P  $\longleftrightarrow ((\lambda x. \ln(\text{of-real } x + s) - \ln(\text{of-real } x)) \longrightarrow 0) \text{ at-top}$ 
using  $s$  by (intro filterlim-cong refl)
      (auto intro: eventually-mono [OF eventually-gt-at-top[of 1::real]])
finally have  $(\lambda x. (1/2) * (\ln(\text{of-real}(real x) + s) - \ln(\text{of-real}(real x)))) \longrightarrow 0$ 
by (rule tendsto-mult-right-zero[OF filterlim-compose[OF - filterlim-real-sequentially]])
thus  $(\lambda x. (1/2) * (\ln(\text{of-nat } x + s) - \ln(\text{of-nat } x))) \longrightarrow 0$  by simp

next
have  $((\lambda x. x * (\ln(1 + s / \text{of-real } x))) \longrightarrow s) \text{ at-top (is ?P)}$ 
by (rule stirling-limit-aux2)
also have  $\ln(1 + s / \text{of-real } x) = \ln(\text{of-real } x + s) - \ln(\text{of-real } x)$  if  $x > 1$ 

```

```

for x
  using that s Ln-divide-of-real [of x of-real x + s, symmetric] neq[of x]
  by (auto simp: Ln-of-real field-simps)
  hence ?P  $\longleftrightarrow$  (( $\lambda x.$  of-real x * (ln (of-real x + s) - ln (of-real x)))  $\longrightarrow$  s)
  at-top
    by (intro filterlim-cong refl)
      (auto intro: eventually-mono[OF eventually-gt-at-top[of 1::real]])
    finally have ( $\lambda n.$  of-real (real n) * (ln (of-real (real n) + s) - ln (of-real (real n))))  $\longrightarrow$  s
      by (rule filterlim-compose[OF - filterlim-real-sequentially])
      thus ( $\lambda n.$  of-nat n * (ln (of-nat n + s) - ln (of-nat n)))  $\longrightarrow$  s by simp
    qed simp-all
    also have ?this  $\longleftrightarrow$  (( $\lambda N.$  integral {0..real N} ( $\lambda x.$  -pbernpoly 1 x / (x + s)))
   $\longrightarrow$ 
    ln-Gamma s + s - (s - 1/2) * ln s - ln (2 * pi) / 2) at-top
  using integral-pbernpoly-1-aux
  by (intro filterlim-cong refl)
    (auto intro: eventually-mono[OF eventually-gt-at-top[of 0::nat]])
  also have ( $\lambda N.$  integral {0..real N} ( $\lambda x.$  -pbernpoly 1 x / (x + s))) =
    ( $\lambda N.$  -integral {0..real N} ( $\lambda x.$  pbernpoly 1 x / (x + s)))
  by (simp add: fun-eq-iff)
  finally show ?thesis by (simp add: tends-to-minus-cancel-left [symmetric] algebra-simps)
qed

```

**qualified lemma** pbernpoly-integral-conv-pbernpoly-integral-Suc:

**assumes**  $n \geq 1$

**shows** integral {0..real N} ( $\lambda x.$  pbernpoly n x / (x + s)  $\wedge$  n) =  
 $\quad$  of-real (pbernpoly (Suc n) (real N)) / (of-nat (Suc n) \* (s + of-nat N)  
 $\wedge$  n) -  
 $\quad$  of-real (bernoulli (Suc n)) / (of-nat (Suc n) \* s  $\wedge$  n) + of-nat n / of-nat  
 $\quad$  (Suc n) \*  
 $\quad$  integral {0..real N} ( $\lambda x.$  of-real (pbernpoly (Suc n) x) / (of-real x +  
 $\quad$  s)  $\wedge$  Suc n)

**proof -**

**note** [derivative-intros] = has-field-derivative-pbernpoly-Suc'

**define** I where  $I = -\text{of-real}(\text{pbernpoly}(\text{Suc } n)(\text{of-nat } N)) / (\text{of-nat}(\text{Suc } n) * (\text{of-nat } N + s) \wedge n) +$   
 $\quad \text{of-real}(\text{bernoulli}(\text{Suc } n) / \text{real}(\text{Suc } n)) / s \wedge n +$   
 $\quad \text{integral}\{0..\text{real } N\}(\lambda x. \text{of-real}(\text{pbernpoly} n x) / (\text{of-real } x + s) \wedge n)$

**have** (( $\lambda x.$  (-of-nat n \* inverse (of-real x + s)  $\wedge$  Suc n) \*  
 $\quad$  (of-real (pbernpoly (Suc n) x) / (of-nat (Suc n))))  
 $\quad$  has-integral -I) {0..real N})

**proof** (rule integration-by-parts-interior-strong[OF bounded-bilinear-mult])

**fix** x :: real **assume** x:  $x \in \{0 < .. < \text{real } N\} - \text{real} ` \{0..N\}$

**have**  $x \notin \mathbb{Z}$

**proof**

**assume**  $x \in \mathbb{Z}$

```

then obtain n where x = of-int n by (auto elim!: Ints-cases)
with x have x': x = of-nat (nat n) by simp
from x show False by (auto simp: x')
qed
hence ((λx. of-real (pbernpoly (Suc n) x) / of-nat (Suc n))) has-vector-derivative
complex-of-real (pbernpoly n x)) (at x)
by (intro has-vector-derivative-of-real) (auto intro!: derivative-eq-intros)
thus ((λx. of-real (pbernpoly (Suc n) x) / of-nat (Suc n))) has-vector-derivative
complex-of-real (pbernpoly n x)) (at x) by simp
from x s have complex-of-real x + s ≠ 0
by (auto simp: complex-eq-iff complex-nonpos-Reals-iff)
thus ((λx. inverse (of-real x + s) ^ n) has-vector-derivative
- of-nat n * inverse (of-real x + s) ^ Suc n) (at x) using x s assms
by (auto intro!: derivative-eq-intros has-vector-derivative-real-field simp: di-
vide-simps power-add [symmetric]
simp del: power-Suc)
next
have complex-of-real x + s ≠ 0 if x ≥ 0 for x
using that s by (auto simp: complex-eq-iff complex-nonpos-Reals-iff)
thus continuous-on {0..real N} (λx. inverse (of-real x + s) ^ n)
continuous-on {0..real N} (λx. complex-of-real (pbernpoly (Suc n) x) /
of-nat (Suc n))
using assms s by (auto intro!: continuous-intros simp del: of-nat-Suc)
next
have ((λx. inverse (of-real x + s) ^ n * of-real (pbernpoly n x)) has-integral
pbernpoly (Suc n) (of-nat N) / (of-nat (Suc n) * (of-nat N + s) ^ n) -
of-real (beroulli (Suc n) / real (Suc n)) / s ^ n - I) {0..real N}
using integrable-ln-Gamma-aux[of n N] assms
by (auto simp: I-def has-integral-integral divide-simps)
thus ((λx. inverse (of-real x + s) ^ n * of-real (pbernpoly n x)) has-integral
inverse (of-real (real N) + s) ^ n * (of-real (pbernpoly (Suc n) (real
N)) /
of-nat (Suc n)) -
inverse (of-real 0 + s) ^ n * (of-real (pbernpoly (Suc n) 0) / of-nat
(Suc n)) - I)
{0..real N} by (simp-all add: field-simps)
qed simp-all
also have (λx. - of-nat n * inverse (of-real x + s) ^ Suc n * (of-real (pbernpoly
(Suc n) x) /
of-nat (Suc n))) =
(λx. - (of-nat n / of-nat (Suc n) * of-real (pbernpoly (Suc n) x) /
(of-real x + s) ^ Suc n))
by (simp add: divide-simps fun-eq-iff)
finally have ((λx. - (of-nat n / of-nat (Suc n) * of-real (pbernpoly (Suc n) x) /
(of-real x + s) ^ Suc n)) has-integral - I) {0..real N} .
from has-integral-neg[OF this] show ?thesis
by (auto simp add: I-def has-integral-iff algebra-simps integral-mult-right [symmetric]
simp del: power-Suc of-nat-Suc )

```

qed

```
lemma pbernpoly-over-power-tendsto-0:
  assumes n > 0
  shows  ( $\lambda x. \text{of-real}(\text{pbernpoly}(\text{Suc } n)(\text{real } x)) / (\text{of-nat}(\text{Suc } n) * (s + \text{of-nat } x) ^ n)$ ) —→ 0
proof -
  from s have neq:  $s + \text{of-nat } n \neq 0$  for n
    by (auto simp: complex-eq-iff complex-nonpos-Reals-iff)
  obtain c where c:  $\bigwedge x. \text{norm}(\text{pbernpoly}(\text{Suc } n) x) \leq c$ 
    using bounded-pbernpoly by auto
  have eventually ( $\lambda x. \text{real } x + \text{Re } s > 0$ ) at-top
    by real-asym
  hence eventually ( $\lambda x. \text{norm}(\text{of-real}(\text{pbernpoly}(\text{Suc } n)(\text{real } x)) / (\text{of-nat}(\text{Suc } n) * (s + \text{of-nat } x) ^ n)) \leq (c / \text{real}(\text{Suc } n)) / (\text{real } x + \text{Re } s) ^ n$ ) at-top
    using eventually-gt-at-top[of 0::nat]
  proof eventually-elim
    case (elim x)
    have norm (of-real (pbernpoly (Suc n) (real x)) /
      (of-nat (Suc n) * (s + of-nat x) ^ n)) ≤
      (c / real (Suc n)) / norm (s + of-nat x) ^ n (is - ≤ ?rhs) using c[of x]
      by (auto simp: norm-divide norm-mult norm-power neq field-simps simp del:
        of-nat-Suc)
    also have ( $\text{real } x + \text{Re } s \leq \text{cmod}(s + \text{of-nat } x)$ )
      using complex-Re-le-cmod[of s + of-nat x] s by (auto simp add: complex-nonpos-Reals-iff)
    hence ?rhs ≤ (c / real (Suc n)) / (real x + Re s) ^ n using s elim c[of 0]
      neq[of x]
      by (intro divide-left-mono power-mono mult-pos-pos divide-nonneg-pos zero-less-power)
    auto
    finally show ?case .
  qed
  moreover have ( $\lambda x. (c / \text{real}(\text{Suc } n)) / (\text{real } x + \text{Re } s) ^ n$ ) —→ 0
    using ‹n > 0› by real-asym
  ultimately show ?thesis by (rule Lim-null-comparison)
qed
```

```
lemma convergent-stirling-integral:
  assumes n > 0
  shows convergent ( $\lambda N. \text{integral}\{0..\text{real } N\} (\lambda x. \text{of-real}(\text{pbernpoly } n x) / (\text{of-real } x + s) ^ n)$ ) (is convergent (?f n))
proof -
  have convergent (?f (Suc n)) for n
  proof (induction n)
    case 0
    thus ?case using integral-pbernpoly-1 by (auto intro!: convergentI)
  next
    case (Suc n)
```

```

have convergent ( $\lambda N. ?f(Suc n) N -$ 
   $of\text{-real}(pbernpoly(Suc(Suc n))(real N)) /$ 
   $(of\text{-nat}(Suc(Suc n)) * (s + of\text{-nat}N) \wedge Suc n) +$ 
   $of\text{-real}(bernioulli(Suc(Suc n)) / (real(Suc(Suc n)))) / s \wedge Suc n)$ 
(is convergent ?g)
by (intro convergent-add convergent-diff Suc
  convergent-const convergentI[OF pbernpoly-over-power-tendsto-0]) simp-all
also have ?g = ( $\lambda N. of\text{-nat}(Suc n) / of\text{-nat}(Suc(Suc n)) * ?f(Suc(Suc n))$ 
N) using s
by (subst pbernpoly-integral-conv-pbernpoly-integral-Suc)
  (auto simp: fun-eq-iff field-simps simp del: of-nat-Suc power-Suc)
also have convergent ...  $\longleftrightarrow$  convergent (?f(Suc(Suc n)))
by (intro convergent-mult-const-iff) (simp-all del: of-nat-Suc)
finally show ?case .
qed
from this[of n - 1] assms show ?thesis by simp
qed

lemma stirling-integral-conv-stirling-integral-Suc:
assumes n > 0
shows stirling-integral n s =
  of-nat n / of-nat(Suc n) * stirling-integral(Suc n) s -
  of-real(bernioulli(Suc n)) / (of-nat(Suc n) * s  $\wedge$  n)

proof -
have ( $\lambda N. of\text{-real}(pbernpoly(Suc n)(real N)) / (of\text{-nat}(Suc n) * (s + of\text{-nat}N) \wedge n) -$ 
   $of\text{-real}(bernioulli(Suc n)) / (real(Suc n) * s \wedge n) +$ 
   $integral\{0..real N\}(\lambda x. of\text{-nat}n / of\text{-nat}(Suc n) *$ 
   $(of\text{-real}(pbernpoly(Suc n)x) / (of\text{-real}x + s) \wedge Suc n)))$ 
 $\longrightarrow 0 - of\text{-real}(bernioulli(Suc n)) / (of\text{-nat}(Suc n) * s \wedge n) +$ 
   $of\text{-nat}n / of\text{-nat}(Suc n) * stirling-integral(Suc n) s$  (is ?f  $\longrightarrow$ 
-)
unfolding stirling-integral-def integral-mult-right
using convergent-stirling-integral[of Suc n] assms s
by (intro tendsto-intros pbernpoly-over-power-tendsto-0)
  (auto simp: convergent-LIMSEQ-iff simp del: of-nat-Suc)
also have ?this  $\longleftrightarrow$  ( $\lambda N. integral\{0..real N\}$ 
   $(\lambda x. of\text{-real}(pbernpoly n x) / (of\text{-real}x + s) \wedge n)) \longrightarrow$ 
   $of\text{-nat}n / of\text{-nat}(Suc n) * stirling-integral(Suc n) s -$ 
   $of\text{-real}(bernioulli(Suc n)) / (of\text{-nat}(Suc n) * s \wedge n)$ 
using eventually-gt-at-top[of 0::nat] pbernpoly-integral-conv-pbernpoly-integral-Suc[of
n]
assms unfolding integral-mult-right
by (intro filterlim-cong refl) (auto elim!: eventually-mono simp del: power-Suc)
finally show ?thesis unfolding stirling-integral-def[of n] by (rule limI)
qed

lemma stirling-integral-1-unfold:
assumes m > 0

```

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shows stirling-integral 1 s = stirling-integral m s / of-nat m -
          ( $\sum_{k=1..<m. \text{of-real}(\text{bernoulli}(\text{Suc } k)) / (\text{of-nat } k * \text{of-nat}(\text{Suc } k) *$ 
           $s^{\wedge k})$ )
proof -
  have stirling-integral 1 s = stirling-integral (Suc m) s / of-nat (Suc m) -
          ( $\sum_{k=1..<\text{Suc } m. \text{of-real}(\text{bernoulli}(\text{Suc } k)) / (\text{of-nat } k * \text{of-nat}(\text{Suc } k)$ 
           $* s^{\wedge k})$ ) for m
  proof (induction m)
    case (Suc m)
      let ?C = ( $\sum_{k=1..<\text{Suc } m. \text{of-real}(\text{bernoulli}(\text{Suc } k)) / (\text{of-nat } k * \text{of-nat}(\text{Suc } k)$ 
       $* s^{\wedge k})$ )
      note Suc.IH
      also have stirling-integral (Suc m) s / of-nat (Suc m) =
          stirling-integral (Suc (Suc m)) s / of-nat (Suc (Suc m)) -
          of-real (bernoulli (Suc (Suc m))) /
          (of-nat (Suc m) * of-nat (Suc (Suc m)) * s^{\wedge Suc m})
      (is - = ?A - ?B) by (subst stirling-integral-conv-stirling-integral-Suc)
          (simp-all del: of-nat-Suc power-Suc add: divide-simps)
      also have ?A - ?B - ?C = ?A - (?B + ?C) by (rule diff-diff-eq)
      also have ?B + ?C = ( $\sum_{k=1..<\text{Suc } (Suc m). \text{of-real}(\text{bernoulli}(\text{Suc } k)) /$ 
       $(\text{of-nat } k * \text{of-nat}(\text{Suc } k) * s^{\wedge k})$ )
      using s by (simp add: divide-simps)
      finally show ?case .
  qed simp-all
  note this[of m - 1]
  also from assms have Suc (m - 1) = m by simp
  finally show ?thesis .
qed

lemma ln-Gamma-stirling-complex:
assumes m > 0
shows ln-Gamma s = (s - 1 / 2) * ln s - s + ln (2 * pi) / 2 +
          ( $\sum_{k=1..<m. \text{of-real}(\text{bernoulli}(\text{Suc } k)) / (\text{of-nat } k * \text{of-nat}(\text{Suc } k)$ 
           $* s^{\wedge k})$ ) -
          stirling-integral m s / of-nat m
proof -
  have ln-Gamma s = (s - 1 / 2) * ln s - s + ln (2 * pi) / 2 - stirling-integral
  1 s
  using limI[OF integral-pbernpoly-1] by (simp add: stirling-integral-def alge-
  bra-simps)
  also have stirling-integral 1 s = stirling-integral m s / of-nat m -
          ( $\sum_{k=1..<m. \text{of-real}(\text{bernoulli}(\text{Suc } k)) / (\text{of-nat } k * \text{of-nat}(\text{Suc } k)$ 
           $* s^{\wedge k})$ )
  using assms by (rule stirling-integral-1-unfold)
  finally show ?thesis by simp
qed

lemma LIMSEQ-stirling-integral:
n > 0  $\implies$  ( $\lambda x. \text{integral} \{0.. \text{real } x\} (\lambda x. \text{of-real}(\text{pbernpoly } n x)) / (\text{of-real } x + s)$ 
```

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 $\wedge n))$ 
 $\longrightarrow \text{stirling-integral } n s \text{ unfolding stirling-integral-def}$ 
using convergent-stirling-integral[of  $n$ ] by (simp only: convergent-LIMSEQ-iff)

end

lemmas has-integral-of-real = has-integral-linear[OF - bounded-linear-of-real, unfolded o-def]
lemmas integral-of-real = integral-linear[OF - bounded-linear-of-real, unfolded o-def]

lemma integrable-ln-Gamma-aux-real:
assumes  $0 < s$ 
shows  $(\lambda x. pbernpoly n x / (x + s)^\wedge n)$  integrable-on {0..real  $N$ }
proof -
  have  $(\lambda x. \text{complex-of-real } (pbernpoly n x / (x + s)^\wedge n))$  integrable-on {0..real  $N$ }
    using integrable-ln-Gamma-aux[of of-real  $s n N$ ] assms by simp
    from integrable-linear[OF this bounded-linear-Re] show ?thesis
      by (simp only: o-def Re-complex-of-real)
qed

lemma
assumes  $x > 0 n > 0$ 
shows stirling-integral-complex-of-real:
  stirling-integral  $n$  (complex-of-real  $x$ ) = of-real (stirling-integral  $n x$ )
and LIMSEQ-stirling-integral-real:
   $(\lambda N. \text{integral } \{0..real N\} (\lambda t. pbernpoly n t / (t + x)^\wedge n))$ 
   $\longrightarrow \text{stirling-integral } n x$ 
and stirling-integral-real-convergent:
  convergent  $(\lambda N. \text{integral } \{0..real N\} (\lambda t. pbernpoly n t / (t + x)^\wedge n))$ 
proof -
  have  $(\lambda N. \text{integral } \{0..real N\} (\lambda t. \text{of-real } (pbernpoly n t / (t + x)^\wedge n)))$ 
     $\longrightarrow \text{stirling-integral } n (\text{complex-of-real } x)$ 
  using LIMSEQ-stirling-integral[of complex-of-real  $x n$ ] assms by simp
  hence  $(\lambda N. \text{of-real } (\text{integral } \{0..real N\} (\lambda t. pbernpoly n t / (t + x)^\wedge n)))$ 
     $\longrightarrow \text{stirling-integral } n (\text{complex-of-real } x)$ 
  using integrable-ln-Gamma-aux-real[OF assms(1), of  $n$ ]
  by (subst (asm) integral-of-real) simp
  from tendsto-Re[OF this]
  have  $(\lambda N. \text{integral } \{0..real N\} (\lambda t. pbernpoly n t / (t + x)^\wedge n))$ 
     $\longrightarrow \text{Re } (\text{stirling-integral } n (\text{complex-of-real } x))$  by simp
  thus convergent  $(\lambda N. \text{integral } \{0..real N\} (\lambda t. pbernpoly n t / (t + x)^\wedge n))$ 
    by (rule convergentI)
  thus  $(\lambda N. \text{integral } \{0..real N\} (\lambda t. pbernpoly n t / (t + x)^\wedge n))$ 
     $\longrightarrow \text{stirling-integral } n x$  unfolding stirling-integral-def
    by (simp add: convergent-LIMSEQ-iff)
  from tendsto-of-real[OF this, where 'a = complex]
    integrable-ln-Gamma-aux-real[OF assms(1), of  $n$ ]
  have  $(\lambda x a. \text{integral } \{0..real xa\})$ 

```

```


$$(\lambda xa. \text{complex-of-real} (\text{pbernpoly } n \text{ } xa) / (\text{complex-of-real} \text{ } xa + x)^n))$$


$$\longrightarrow \text{complex-of-real} (\text{stirling-integral} \text{ } n \text{ } x)$$

by (subst (asm) integral-of-real [symmetric]) simp-all
from LIMSEQ-unique[OF this LIMSEQ-stirling-integral[of complex-of-real x n]]
assms
show stirling-integral n (complex-of-real x) = of-real (stirling-integral n x) by
simp
qed

lemma ln-Gamma-stirling-real:
assumes x > (0 :: real) m > (0::nat)
shows ln-Gamma x = (x - 1 / 2) * ln x - x + ln (2 * pi) / 2 +

$$(\sum k = 1..<m. \text{bernoulli} (\text{Suc } k) / (\text{of-nat } k * \text{of-nat} (\text{Suc } k) * x^k))$$

-
stirling-integral m x / of-nat m
proof -
from assms have complex-of-real (ln-Gamma x) = ln-Gamma (complex-of-real x)
by (simp add: ln-Gamma-complex-of-real)
also have ln-Gamma (complex-of-real x) = complex-of-real (

$$(x - 1 / 2) * \ln x - x + \ln (2 * \pi) / 2 +$$


$$(\sum k = 1..<m. \text{bernoulli} (\text{Suc } k) / (\text{of-nat } k * \text{of-nat} (\text{Suc } k) * x^k))$$

)
stirling-integral m x / of-nat m using assms
by (subst ln-Gamma-stirling-complex[of - m])
(basic-all add: Ln-of-real stirling-integral-complex-of-real)
finally show ?thesis by (subst (asm) of-real-eq-iff)
qed

lemma stirling-integral-bound-aux:
assumes n: n > (1::nat)
obtains c where  $\bigwedge s. \text{Re } s > 0 \implies \text{norm} (\text{stirling-integral } n \text{ } s) \leq c / \text{Re } s^{n-1}$ 
proof -
obtain c where c: norm (pbernpoly n x) ≤ c for x by (rule bounded-pbernpoly[of n])
blast
have c': pbernpoly n x ≤ c for x using c[of x] by (simp add: abs-real-def split: if-splits)
from c[of 0] have c-nonneg: c ≥ 0 by simp
have norm (stirling-integral n s) ≤ c / (real n - 1) / Re s^(n-1) if s: Re s > 0 for s
proof (rule Lim-norm-ubound[OF - LIMSEQ-stirling-integral])
have pos: x + norm s > 0 if x ≥ 0 for x using s that by (intro add-nonneg-pos)
auto
have nz: of-real x + s ≠ 0 if x ≥ 0 for x using s that by (auto simp: complex-eq-iff)
let ?bound = λN. c / (Re s^(n-1) * (real n - 1)) -

```

```

c / ((real N + Re s) ^ (n - 1) * (real n - 1))
show eventually (λN. norm (integral {0..real N}
  (λx. of-real (pbernpoly n x) / (of-real x + s) ^ n)) ≤
  c / (real n - 1) / Re s ^ (n - 1)) at-top
using eventually-gt-at-top[of 0::nat]
proof eventually-elim
  case (elim N)
    let ?F = λx. -c / ((x + Re s) ^ (n - 1) * (real n - 1))
    from n s have ((λx. c / (x + Re s) ^ n) has-integral (?F (real N) - ?F 0))
  {0..real N}
    by (intro fundamental-theorem-of-calculus)
      (auto intro!: derivative-eq-intros simp: divide-simps power-diff add-eq-0-iff2
        has-real-derivative-iff-has-vector-derivative [symmetric])
    also have ?F (real N) - ?F 0 = ?bound N by simp
    finally have *: ((λx. c / (x + Re s) ^ n) has-integral ?bound N) {0..real N}

    .
    have norm (integral {0..real N} (λx. of-real (pbernpoly n x) / (of-real x + s)
  ^ n)) ≤
      integral {0..real N} (λx. c / (x + Re s) ^ n)
    proof (intro integral-norm-bound-integral integrable-ln-Gamma-aux s ballI)
      fix x assume x: x ∈ {0..real N}
      have norm (of-real (pbernpoly n x) / (of-real x + s) ^ n) ≤ c / norm (of-real
  x + s) ^ n
      unfolding norm-divide norm-power using c by (intro divide-right-mono)
  simp-all
      also have ... ≤ c / (x + Re s) ^ n
      using x c c-nonneg s nz[of x] complex-Re-le-cmod[of of-real x + s]
        by (intro divide-left-mono power-mono mult-pos-pos zero-less-power
  add-nonneg-pos) auto
      finally show norm (of-real (pbernpoly n x) / (of-real x + s) ^ n) ≤ ... .
    qed (insert n s * pos nz c, auto simp: complex-nonpos-Reals-iff)
    also have ... = ?bound N using * by (simp add: has-integral-iff)
    also have ... ≤ c / (Re s ^ (n - 1) * (real n - 1)) using c-nonneg elim s
  n by simp
    also have ... = c / (real n - 1) / (Re s ^ (n - 1)) by simp
    finally show norm (integral {0..real N} (λx. of-real (pbernpoly n x) /
  (of-real x + s) ^ n)) ≤ c / (real n - 1) / Re s ^ (n - 1) .

  qed
  qed (insert s n, simp-all add: complex-nonpos-Reals-iff)
  thus ?thesis by (rule that)
qed

lemma stirling-integral-bound-aux-integral1:
  fixes a b c :: real and n :: nat
  assumes a ≥ 0 b > 0 c ≥ 0 n > 1 l < a - b r > a + b
  shows ((λx. c / max b |x - a| ^ n) has-integral
  2*c*(n / (n - 1))/b^(n-1) - c/(n-1) * (1/(a-l)^(n-1) + 1/(r-a)^(n-1)))
  {l..r}
proof -

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```

define x1 x2 where x1 = a - b and x2 = a + b
define F1 where F1 = ( $\lambda x::real. c / (a - x) \wedge (n - 1) / (n - 1)$ )
define F3 where F3 = ( $\lambda x::real. -c / (x - a) \wedge (n - 1) / (n - 1)$ )
have deriv: (F1 has-vector-derivative ( $c / (a - x) \wedge n$ )) (at x within A)
      (F3 has-vector-derivative ( $c / (x - a) \wedge n$ )) (at x within A)
  if  $x \neq a$  for  $x :: real$  and A
  unfolding F1-def F3-def using assms that
  by (auto intro!: derivative-eq-intros simp: divide-simps power-diff add-eq-0-iff2
       simp flip: has-real-derivative-iff-has-vector-derivative)

from assms have (( $\lambda x. c / (a - x) \wedge n$ ) has-integral (F1 x1 - F1 l)) {l..x1}
  by (intro fundamental-theorem-of-calculus deriv) (auto simp: x1-def max-def
split: if-splits)
also have ?this  $\longleftrightarrow$  (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral (F1 x1 - F1 l))
{l..x1}
  using assms
  by (intro has-integral-spike-finite-eq[of {l}]) (auto simp: x1-def max-def split:
if-splits)
finally have I1: (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral (F1 x1 - F1 l))
{l..x1} .

have (( $\lambda x. c / b \wedge n$ ) has-integral (x2 - x1) * c / b  $\wedge n$ ) {x1..x2}
  using has-integral-const-real[of c / b  $\wedge n$  x1 x2] assms by (simp add: x1-def
x2-def)
also have ?this  $\longleftrightarrow$  (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral ((x2 - x1) * c /
b  $\wedge n$ ) {x1..x2})
  by (intro has-integral-spike-finite-eq[of {x1, x2}])
    (auto simp: x1-def x2-def split: if-splits)
finally have I2: (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral ((x2 - x1) * c / b  $\wedge$ 
n)) {x1..x2} .

from assms have I3: (( $\lambda x. c / (x - a) \wedge n$ ) has-integral (F3 r - F3 x2)) {x2..r}
  by (intro fundamental-theorem-of-calculus deriv) (auto simp: x2-def min-def
split: if-splits)
also have ?this  $\longleftrightarrow$  (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral (F3 r - F3 x2))
{x2..r}
  using assms
  by (intro has-integral-spike-finite-eq[of {r}]) (auto simp: x2-def min-def split:
if-splits)
finally have I3: (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral (F3 r - F3 x2))
{x2..r} .

have (( $\lambda x. c / \max b |x - a| \wedge n$ ) has-integral (F1 x1 - F1 l) + ((x2 - x1) *
c / b  $\wedge n$ ) + (F3 r - F3 x2)) {l..r}
  using assms
  by (intro has-integral-combine[OF - - has-integral-combine[OF - - I1 I2] I3])
    (auto simp: x1-def x2-def)
also have (F1 x1 - F1 l) + ((x2 - x1) * c / b  $\wedge n$ ) + (F3 r - F3 x2) =
      F1 x1 - F1 l + F3 r - F3 x2 + (x2 - x1) * c / b  $\wedge n$ 

```

```

by (simp add: algebra-simps)
also have  $x_2 - x_1 = 2 * b$ 
  using assms by (simp add: x2-def x1-def min-def max-def)
also have  $2 * b * c / b^{\wedge} n = 2 * c / b^{\wedge} (n - 1)$ 
  using assms by (simp add: power-diff field-simps)
also have  $F_1 x_1 - F_1 l + F_3 r - F_3 x_2 =$ 
 $c/(n-1) * (2/b^{\wedge}(n-1) - 1/(a-l)^{\wedge}(n-1) - 1/(r-a)^{\wedge}(n-1))$ 
  using assms by (simp add: x1-def x2-def F1-def F3-def field-simps del: of-nat-diff)
also have ... +  $2 * c / b^{\wedge} (n - 1) =$ 
 $2*c*(1 + 1/(n-1))/b^{\wedge}(n-1) - c/(n-1) * (1/(a-l)^{\wedge}(n-1) +$ 
 $1/(r-a)^{\wedge}(n-1))$ 
  using assms by (simp add: field-simps del: of-nat-diff)
also have  $1 + 1 / (n - 1) = n / (n - 1)$ 
  using assms by (simp add: field-simps)
finally show ?thesis .
qed

```

```

lemma stirling-integral-bound-aux-integral2:
fixes a b c :: real and n :: nat
assumes a ≥ 0 b > 0 c ≥ 0 n > 1
obtains I where ((λx. c / max b |x - a| ^ n) has-integral I) {l..r}
I ≤ 2 * c * (n / (n - 1)) / b ^ (n-1)

proof -
define l' where l' = min l (a - b - 1)
define r' where r' = max r (a + b + 1)

define A where A = 2 * c * (n / (n - 1)) / b ^ (n - 1)
define B where B = c / real (n - 1) * (1 / (a - l') ^ (n - 1) + 1 / (r' -
a) ^ (n - 1))

have has-int: ((λx. c / max b |x - a| ^ n) has-integral (A - B)) {l'..r'}
  using assms unfolding A-def B-def
  by (intro stirling-integral-bound-aux-integral1) (auto simp: l'-def r'-def)
have (λx. c / max b |x - a| ^ n) integrable-on {l..r}
  by (rule integrable-on-subinterval[OF has-integral-integrable[OF has-int]])
  (auto simp: l'-def r'-def)
then obtain I where has-int': ((λx. c / max b |x - a| ^ n) has-integral I) {l..r}
  by (auto simp: integrable-on-def)

from assms have I ≤ A - B
  by (intro has-integral-subset-le[OF - has-int' has-int]) (auto simp: l'-def r'-def)
also have ... ≤ A
  using assms by (simp add: B-def l'-def r'-def)
finally show ?thesis using that[of I] has-int' unfolding A-def by blast
qed

lemma stirling-integral-bound-aux':
assumes n: n > (1::nat) and α: α ∈ {0 <.. < pi}
obtains c where ⋀s::complex. s ∈ complex-cone' α - {0} ==>

```

```

norm (stirling-integral n s) ≤ c / norm s ^ (n - 1)

proof –
  obtain c where c: norm (pbernpoly n x) ≤ c for x by (rule bounded-pbernpoly[of n]) blast
    have c': pbernpoly n x ≤ c for x using c[of x] by (simp add: abs-real-def split: if-splits)
    from c[of 0] have c-nonneg: c ≥ 0 by simp

  define D where D = c * Beta (- (real-of-int (- int n) / 2) - 1 / 2) (1 / 2)
  / 2
  define C where C = max D (2*c*(n/(n-1))/sin α^(n-1))

  have *: norm (stirling-integral n s) ≤ C / norm s ^ (n - 1)
    if s: s ∈ complex-cone' α - {0} for s :: complex
    proof (rule Lim-norm-ubound[OF - LIMSEQ-stirling-integral])
      from s α have Arg: |Arg s| ≤ α by (auto simp: complex-cone-altdef)
      have s': s ∉ ℝ≤₀
        using complex-cone-inter-nonpos-Reals[of -α α] α s by auto
      from s have [simp]: s ≠ 0 by auto

    show eventually (λN. norm (integral {0..real N}
      (λx. of-real (pbernpoly n x) / (of-real x + s) ^ n)) ≤
      C / norm s ^ (n - 1)) at-top
      using eventually-gt-at-top[of 0::nat]
    proof eventually-elim
      case (elim N)
      show ?case
      proof (cases Re s > 0)
        case True
        have int: ((λx. c * (x^2 + norm s^2) powr (-n / 2)) has-integral
          D * (norm s ^ 2) powr (-n / 2 + 1 / 2)) {0<..}
          using has-integral-mult-left[OF has-integral-Beta3[of -n/2 norm s ^ 2],
          of c] assms True
          unfolding D-def by (simp add: algebra-simps)
          hence int': ((λx. c * (x^2 + norm s^2) powr (-n / 2)) has-integral
            D * (norm s ^ 2) powr (-n / 2 + 1 / 2)) {0..}
            by (subst has-integral-interior [symmetric]) simp-all
          hence integrable: (λx. c * (x^2 + norm s^2) powr (-n / 2)) integrable-on
          {0..}
          by (simp add: has-integral-iff)

        have norm (integral {0..real N} (λx. of-real (pbernpoly n x) / (of-real x +
          s) ^ n)) ≤
          integral {0..real N} (λx. c * (x^2 + norm s^2) powr (-n / 2))
        proof (intro integral-norm-bound-integral s ballI integrable-ln-Gamma-aux)
          have [simp]: {0<..} - {0::real..} = {} {0..} - {0<..} = {0::real}
          by auto
          have (λx. c * (x^2 + (cmod s)^2) powr (real-of-int (- int n) / 2)) integrable-on
          {0<..}

```

```

    using int by (simp add: has-integral-iff)
    also have ?this  $\longleftrightarrow$   $(\lambda x. c * (x^2 + (cmod s)^2)) \text{ powr} (\text{real-of-int } (- \text{int } n) / 2)$ 
    integrable-on {0..}
        by (intro integrable-spoke-set-eq) auto
        finally show  $(\lambda x. c * (x^2 + (cmod s)^2)) \text{ powr} (\text{real-of-int } (- \text{int } n) / 2)$ 
    integrable-on
        {0..real N} by (rule integrable-on-subinterval) auto
    next
        fix x assume x:  $x \in \{0..real N\}$ 
        have nz: complex-of-real x + s  $\neq 0$ 
            using True x by (auto simp: complex-eq-iff)
            have norm (of-real (pbernpoly n x) / (of-real x + s) ^ n)  $\leq c / \text{norm}$ 
            (of-real x + s) ^ n
                unfolding norm-divide norm-power using c by (intro divide-right-mono)
            simp-all
            also have ...  $\leq c / \sqrt{x^2 + \text{norm } s^2}^n$ 
            proof (intro divide-left-mono mult-pos-pos zero-less-power power-mono)
                show  $\sqrt{x^2 + (cmod s)^2} \leq cmod (\text{complex-of-real } x + s)$ 
                    using x True by (simp add: cmod-def algebra-simps power2-eq-square)
                qed (use x True c-nonneg assms nz in (auto simp: add-nonneg-pos))
                also have  $\sqrt{(x^2 + \text{norm } s^2)^n} = (x^2 + \text{norm } s^2)^{\text{powr} (1/2 * n)}$ 
                    by (subst powr-powr [symmetric], subst powr-realpow)
                    (auto simp: powr-half-sqrt add-nonneg-pos)
                also have  $c / \dots = c * (x^2 + \text{norm } s^2)^{\text{powr} (-n / 2)}$ 
                    by (simp add: powr-minus field-simps)
                finally show  $\text{norm} (\text{complex-of-real } (pbernpoly n x) / (\text{complex-of-real } x + s)^n) \leq \dots$ 
            qed fact+
            also have ...  $\leq \text{integral } \{0..\} (\lambda x. c * (x^2 + \text{norm } s^2)^{\text{powr} (-n / 2)})$ 
                using c-nonneg
                by (intro integral-subset-le integrable integrable-on-subinterval[OF integrable]) auto
            also have ...  $= D * (\text{norm } s^2)^{\text{powr} (-n / 2 + 1 / 2)}$ 
                using int' by (simp add: has-integral-iff)
            also have  $(\text{norm } s^2)^{\text{powr} (-n / 2 + 1 / 2)} = \text{norm } s \text{ powr} (2 * (-n / 2 + 1 / 2))$ 
                by (subst powr-powr [symmetric]) auto
            also have ...  $= \text{norm } s \text{ powr} (-\text{real } (n - 1))$ 
                using assms by (simp add: of-nat-diff)
            also have  $D * \dots = D / \text{norm } s^{(n - 1)}$ 
                by (auto simp: powr-minus powr-realpow field-simps)
            also have ...  $\leq C / \text{norm } s^{(n - 1)}$ 
                by (intro divide-right-mono) (auto simp: C-def)
            finally show  $\text{norm} (\text{integral } \{0..real N\} (\lambda x. \text{of-real } (pbernpoly n x) / (\text{of-real } x + s)^n)) \leq \dots$ .
        next

```

```

case False
have cos |Arg s| = cos (Arg s)
  by (simp add: abs-if)
also have cos (Arg s) = Re (rcis (norm s) (Arg s)) / norm s
  by (subst Re-rcis) auto
also have ... = Re s / norm s
  by (subst rcis-cmod-Arg) auto
also have ... ≤ cos (pi / 2)
  using False by (auto simp: field-simps)
finally have |Arg s| ≥ pi / 2
  using Arg α by (subst (asm) cos-mono-le-eq) auto

have sin α * norm s = sin (pi - α) * norm s
  by simp
also have ... ≤ sin (pi - |Arg s|) * norm s
  using α Arg ⟨|Arg s| ≥ pi / 2⟩
  by (intro mult-right-mono sin-monotone-2pi-le) auto
also have sin |Arg s| ≥ 0
  using Arg-bounded[of s] by (intro sin-ge-zero) auto
hence sin (pi - |Arg s|) = |sin |Arg s|| 
  by simp
also have ... = |sin (Arg s)|
  by (simp add: abs-if)
also have ... * norm s = |Im (rcis (norm s) (Arg s))|
  by (simp add: abs-mult)
also have ... = |Im s|
  by (subst rcis-cmod-Arg) auto
finally have abs-Im-ge: |Im s| ≥ sin α * norm s .

have [simp]: Im s ≠ 0 s ≠ 0
  using s ⟨s ∉ ℝ≤₀⟩ False
  by (auto simp: cmod-def zero-le-mult-iff complex-nonpos-Reals-iff)
have sin α > 0
  using assms by (intro sin-gt-zero) auto

obtain I where I: ((λx. c / max |Im s| |x + Re s| ^ n) has-integral I)
{0..real N}
  I ≤ 2*c*(n/(n-1)) / |Im s|^(n-1)
  using s c-nonneg assms False
    stirling-integral-bound-aux-integral2[of -Re s |Im s| c n 0 real N] by
  auto

have norm (integral {0..real N} (λx. of-real (pbernpoly n x) / (of-real x +
s) ^ n)) ≤
  integral {0..real N} (λx. c / max |Im s| |x + Re s| ^ n)
proof (intro integral-norm-bound-integral integrable-ln-Gamma-aux s ballI)
  show (λx. c / max |Im s| |x + Re s| ^ n) integrable-on {0..real N}
    using I(1) by (simp add: has-integral-iff)
next

```

```

fix x assume x:  $x \in \{0..real\ N\}$ 
have nz: complex-of-real x + s ≠ 0
  by (auto simp: complex-eq-iff)
have norm (complex-of-real (pbernpoly n x)) / (complex-of-real x + s) ^ n
≤
  c / norm (complex-of-real x + s) ^ n
  unfolding norm-divide norm-power using c[of x] by (intro divide-right-mono) simp-all
  also have ... ≤ c / max |Im s| |x + Re s| ^ n
    using c-nonneg nz abs-Re-le-cmod[of of-real x + s] abs-Im-le-cmod[of of-real x + s]
    by (intro divide-left-mono power-mono mult-pos-pos zero-less-power)
      (auto simp: less-max-iff-disj)
  finally show norm (complex-of-real (pbernpoly n x)) / (complex-of-real x + s) ^ n ≤ ...
  qed (auto simp: complex-nonpos-Real-iff)
  also have ... ≤ 2*c*(n/(n-1)) / |Im s| ^ (n - 1)
    using I by (simp add: has-integral-iff)
  also have ... ≤ 2*c*(n/(n-1)) / (sin α * norm s) ^ (n - 1)
    using ⟨sin α > 0⟩ s c-nonneg abs-Im-ge
    by (intro divide-left-mono mult-pos-pos zero-less-power power-mono
mult-nonneg-nonneg) auto
  also have ... = 2*c*(n/(n-1))/sin α ^ (n-1) / norm s ^ (n - 1)
    by (simp add: field-simps)
  also have ... ≤ C / norm s ^ (n - 1)
    by (intro divide-right-mono) (auto simp: C-def)
  finally show ?thesis .
qed
qed
qed (use that assms complex-cone-inter-nonpos-Real[of -α α] α in auto)
thus ?thesis by (rule that)
qed

lemma stirling-integral-bound:
assumes n > 0
obtains c where
  ∫s. Re s > 0 ==> norm (stirling-integral n s) ≤ c / Re s ^ n
proof -
  let ?f = λs. of-nat n / of-nat (Suc n) * stirling-integral (Suc n) s -
    of-real (bernoulli (Suc n)) / (of-nat (Suc n) * s ^ n)
  from stirling-integral-bound-aux[of Suc n] assms obtain c where
    c: ∫s. Re s > 0 ==> norm (stirling-integral (Suc n) s) ≤ c / Re s ^ n by auto
  define c1 where c1 = real n / real (Suc n) * c
  define c2 where c2 = |bernoulli (Suc n)| / real (Suc n)
  have c2-nonneg: c2 ≥ 0 by (simp add: c2-def)
  show ?thesis
  proof (rule that)
    fix s :: complex assume s: Re s > 0
    hence s': s ∈ ℝ<0 by (auto simp: complex-nonpos-Real-iff)

```

```

have stirling-integral n s = ?f s using s' assms
  by (rule stirling-integral-conv-stirling-integral-Suc)
also have norm ... ≤ norm (of-nat n / of-nat (Suc n) * stirling-integral (Suc
n) s) +
  norm (of-real (bernpoulli (Suc n)) / (of-nat (Suc n) * s ^ n))
  by (rule norm-triangle-ineq4)
also have ... = real n / real (Suc n) * norm (stirling-integral (Suc n) s) +
  c2 / norm s ^ n (is - = ?A + ?B)
  by (simp add: norm-divide norm-mult norm-power c2-def field-simps del:
of-nat-Suc)
also have ?A ≤ real n / real (Suc n) * (c / Re s ^ n)
  by (intro mult-left-mono c s) simp-all
also have ... = c1 / Re s ^ n by (simp add: c1-def)
also have c2 / norm s ^ n ≤ c2 / Re s ^ n using s c2-nonneg
  by (intro divide-left-mono power-mono complex-Re-le-cmod mult-pos-pos
zero-less-power) auto
also have c1 / Re s ^ n + c2 / Re s ^ n = (c1 + c2) / Re s ^ n
  using s by (simp add: field-simps)
finally show norm (stirling-integral n s) ≤ (c1 + c2) / Re s ^ n by – simp-all
qed
qed

lemma stirling-integral-bound':
assumes n > 0 and α ∈ {0<.. $\pi$ }
obtains c where
   $\bigwedge_{s:\text{complex}. s \in \text{complex-cone}' \alpha - \{0\}} \Rightarrow \text{norm}(\text{stirling-integral } n s) \leq c / \text{norm } s ^ n$ 
proof –
let ?f = λs. of-nat n / of-nat (Suc n) * stirling-integral (Suc n) s -
  of-real (bernpoulli (Suc n)) / (of-nat (Suc n) * s ^ n)
from stirling-integral-bound-aux'[of Suc n] assms obtain c where
  c:  $\bigwedge_{s:\text{complex}. s \in \text{complex-cone}' \alpha - \{0\}} \Rightarrow$ 
    norm (stirling-integral (Suc n) s) ≤ c / norm s ^ n by auto
define c1 where c1 = real n / real (Suc n) * c
define c2 where c2 = |bernpoulli (Suc n)| / real (Suc n)
have c2-nonneg: c2 ≥ 0 by (simp add: c2-def)
show ?thesis
proof (rule that)
fix s :: complex assume s: s ∈ complex-cone' α - {0}
have s': s ∉ ℝ≤0
  using complex-cone-inter-nonpos-Reals[of -α α] assms s by auto

have stirling-integral n s = ?f s using s' assms
  by (intro stirling-integral-conv-stirling-integral-Suc) auto
also have norm ... ≤ norm (of-nat n / of-nat (Suc n) * stirling-integral (Suc
n) s) +
  norm (of-real (bernpoulli (Suc n)) / (of-nat (Suc n) * s ^ n))
  by (rule norm-triangle-ineq4)
also have ... = real n / real (Suc n) * norm (stirling-integral (Suc n) s) +

```

```

c2 / norm s ^ n (is - = ?A + ?B)
by (simp add: norm-divide norm-mult norm-power c2-def field-simps del:
of-nat-Suc)
also have ?A ≤ real n / real (Suc n) * (c / norm s ^ n)
by (intro mult-left-mono c s) simp-all
also have ... = c1 / norm s ^ n by (simp add: c1-def)
also have c1 / norm s ^ n + c2 / norm s ^ n = (c1 + c2) / norm s ^ n
using s by (simp add: divide-simps)
finally show norm (stirling-integral n s) ≤ (c1 + c2) / norm s ^ n by -
simp-all
qed
qed

```

**lemma** *stirling-integral-holomorphic* [*holomorphic-intros*]:  
**assumes**  $m: m > 0$  **and**  $A \cap \mathbb{R}_{\leq 0} = \{\}$   
**shows** *stirling-integral m holomorphic-on A*  
**proof** –  
**from** *assms* **have** [simp]:  $z \notin \mathbb{R}_{\leq 0}$  **if**  $z \in A$  **for**  $z$   
**using** *that* **by** auto  
**let**  $?f = \lambda s::complex. of-nat m * ((s - 1 / 2) * Ln s - s + of-real (\ln(2 * pi) / 2) + (\sum_{k=1..<m. of-real (beroulli (Suc k)) / (of-nat k * of-nat (Suc k) * s ^ k)) - ln-Gamma s)$   
**have**  $?f$  **holomorphic-on** *A* **using** *assms*  
**by** (auto intro!: holomorphic-intros simp del: of-nat-Suc elim!: nonpos-Reals-cases)  
**also have**  $?this \longleftrightarrow \text{stirling-integral } m \text{ holomorphic-on } A$   
**using** *assms* **by** (intro holomorphic-cong refl)  
(simp-all add: field-simps ln-Gamma-stirling-complex)  
**finally show** *stirling-integral m holomorphic-on A* .  
qed

**lemma** *stirling-integral-continuous-on-complex* [*continuous-intros*]:  
**assumes**  $m: m > 0$  **and**  $A \cap \mathbb{R}_{\leq 0} = \{\}$   
**shows** *continuous-on A (stirling-integral m :: - ⇒ complex)*  
**by** (intro holomorphic-on-imp-continuous-on stirling-integral-holomorphic *assms*)

**lemma** *has-field-derivative-stirling-integral-complex*:  
**fixes**  $x :: complex$   
**assumes**  $x \notin \mathbb{R}_{\leq 0}$   $n > 0$   
**shows** *(stirling-integral n has-field-derivative deriv (stirling-integral n) x) (at x)*  
**using** *assms*  
**by** (intro holomorphic-derivI[OF *stirling-integral-holomorphic*, of n - $\mathbb{R}_{\leq 0}$ ]) auto

**lemma**

```

assumes n: n > 0 and x > 0
shows deriv-stirling-integral-complex-of-real:
  (deriv  $\widehat{\wedge}$  j) (stirling-integral n) (complex-of-real x) =
    complex-of-real ((deriv  $\widehat{\wedge}$  j) (stirling-integral n) x) (is ?lhs x = ?rhs x)
and differentiable-stirling-integral-real:
  (deriv  $\widehat{\wedge}$  j) (stirling-integral n) field-differentiable at x (is ?thesis2)
proof -
  let ?A = {s. Re s > 0}
  let ?f =  $\lambda j x. (\text{deriv } \widehat{\wedge} j) (\text{stirling-integral } n) (\text{complex-of-real } x)$ 
  let ?f' =  $\lambda j x. \text{complex-of-real} ((\text{deriv } \widehat{\wedge} j) (\text{stirling-integral } n) x)$ 
have [simp]: open ?A by (simp add: open-halvespace-Re-gt)

have ?lhs x = ?rhs x  $\wedge$  (deriv  $\widehat{\wedge}$  j) (stirling-integral n) field-differentiable at x
  if x > 0 for x using that
proof (induction j arbitrary: x)
  case 0
  have (( $\lambda x. \text{Re} (\text{stirling-integral } n (\text{of-real } x))$ ) has-field-derivative
     $\text{Re} (\text{deriv} (\lambda x. \text{stirling-integral } n x) (\text{of-real } x))$  (at x) using 0 n
    by (auto intro!: derivative-intros has-vector-derivative-real-field
      field-differentiable-derivI holomorphic-on-imp-differentiable-at[of - ?A]
      stirling-integral-holomorphic simp: complex-nonpos-Reals-iff)
  also have ?this  $\longleftrightarrow$  (stirling-integral n has-field-derivative
     $\text{Re} (\text{deriv} (\lambda x. \text{stirling-integral } n x) (\text{of-real } x))$  (at x)
    using eventually-nhds-in-open[of {0<..} x] 0 n
    by (intro has-field-derivative-cong-ev refl)
      (auto elim!: eventually-mono simp: stirling-integral-complex-of-real)
  finally have stirling-integral n field-differentiable at x
    by (auto simp: field-differentiable-def)
    with 0 n show ?case by (auto simp: stirling-integral-complex-of-real)
next
  case (Suc j x)
  note IH = conjunct1[OF Suc.IH] conjunct2[OF Suc.IH]
  have *: (deriv  $\widehat{\wedge}$  Suc j) (stirling-integral n) (complex-of-real x) =
    of-real ((deriv  $\widehat{\wedge}$  Suc j) (stirling-integral n) x) if x: x > 0 for x
  proof -
    have deriv ((deriv  $\widehat{\wedge}$  j) (stirling-integral n)) (complex-of-real x) =
      vector-derivative ( $\lambda x. (\text{deriv } \widehat{\wedge} j) (\text{stirling-integral } n) (\text{of-real } x)$ ) (at x)
      using n x
    by (intro vector-derivative-of-real-right [symmetric]
      holomorphic-on-imp-differentiable-at[of - ?A] holomorphic-higher-deriv
      stirling-integral-holomorphic) (auto simp: complex-nonpos-Reals-iff)
    also have ... = vector-derivative ( $\lambda x. \text{of-real} ((\text{deriv } \widehat{\wedge} j) (\text{stirling-integral } n) x)$ ) (at x)
      using eventually-nhds-in-open[of {0<..} x] x
      by (intro vector-derivative-cong-eq) (auto elim!: eventually-mono simp:
        IH(1))
    also have ... = of-real (deriv ((deriv  $\widehat{\wedge}$  j) (stirling-integral n)) x)
      by (intro vector-derivative-of-real-left holomorphic-on-imp-differentiable-at[of

```

```

- ?A]
  field-differentiable-imp-differentiable IH(2) x)
  finally show ?thesis by simp
qed
have ((λx. Re ((deriv ^ Suc j) (stirling-integral n) (of-real x))) has-field-derivative

  Re (deriv ((deriv ^ Suc j) (stirling-integral n)) (of-real x))) (at x)
  using Suc.prems n
  by (intro derivative-intros has-vector-derivative-real-field field-differentiable-derivI
    holomorphic-on-imp-differentiable-at[of - ?A] stirling-integral-holomorphic
    holomorphic-higher-deriv) (auto simp: complex-nonpos-Reals-iff)
also have ?this ↔ ((deriv ^ Suc j) (stirling-integral n) has-field-derivative
  Re (deriv ((deriv ^ Suc j) (stirling-integral n)) (of-real x))) (at x)
  using eventually-nhds-in-open[of {0<..} x] Suc.prems *
  by (intro has-field-derivative-cong-ev refl) (auto elim!: eventually-mono)
finally have (deriv ^ Suc j) (stirling-integral n) field-differentiable at x
  by (auto simp: field-differentiable-def)
  with *[OF Suc.prems] show ?case by blast
qed
from this[OF assms(2)] show ?lhs x = ?rhs x ?thesis2 by blast+
qed

```

Unfortunately, asymptotic power series cannot, in general, be differentiated. However, since *ln-Gamma* is holomorphic on the entire positive real half-space, we can differentiate its asymptotic expansion after all.

To do this, we use an ad-hoc version of the more general approach outlined in Erdelyi's "Asymptotic Expansions" for holomorphic functions: We bound the value of the  $j$ -th derivative of the remainder term at some value  $x$  by applying Cauchy's integral formula along a circle centred at  $x$  with radius  $\frac{1}{2}x$ .

```

lemma deriv-stirling-integral-real-bound:
assumes m: m > 0
shows (deriv ^ j) (stirling-integral m) ∈ O(λx::real. 1 / x ^ (m + j))
proof -
  obtain c where c: ∀s. 0 < Re s ==> cmod (stirling-integral m s) ≤ c / Re s ^ m
    using stirling-integral-bound[OF m] by auto
  have 0 ≤ cmod (stirling-integral m 1) by simp
  also have ... ≤ c using c[of 1] by simp
  finally have c-nonneg: c ≥ 0 .
  define B where B = c * 2 ^ (m + Suc j)
  define B' where B' = B * fact j / 2

  have eventually (λx::real. norm ((deriv ^ j) (stirling-integral m) x) ≤
    B' * norm (1 / x ^ (m + j))) at-top
    using eventually-gt-at-top[of 0::real]
  proof eventually-elim
    case (elim x)

```

```

have  $s \notin \mathbb{R}_{\leq 0}$  if  $s \in cball (\text{of-real } x) (x/2)$  for  $s :: \text{complex}$ 
proof -
  have  $x - Re s \leq \text{norm} (\text{of-real } x - s)$  using complex-Re-le-cmod[of of-real x - s] by simp
    also from that have ...  $\leq x/2$  by (simp add: dist-complex-def)
    finally show ?thesis using elim by (auto simp: complex-nonpos-Reals-iff)
  qed
  hence  $((\lambda u. \text{stirling-integral } m u / (u - \text{of-real } x) \wedge \text{Suc } j) \text{ has-contour-integral}$ 
         $\text{complex-of-real } (2 * pi) * i / \text{fact } j * (deriv \wedge j) (\text{stirling-integral } m) (\text{of-real } x) (\text{circlepath } (\text{of-real } x) (x/2))$ 
        using m elim
        by (intro Cauchy-has-contour-integral-higher-derivative-circlepath
              stirling-integral-continuous-on-complex stirling-integral-holomorphic)
  auto
  hence  $\text{norm} (\text{of-real } (2 * pi) * i / \text{fact } j * (deriv \wedge j) (\text{stirling-integral } m) (\text{of-real } x)) \leq$ 
         $B / x \wedge (m + \text{Suc } j) * (2 * pi * (x / 2))$ 
  proof (rule has-contour-integral-bound-circlepath)
    fix  $u :: \text{complex}$  assume  $\text{dist}: \text{norm} (u - \text{of-real } x) = x / 2$ 
    have  $Re (\text{of-real } x - u) \leq \text{norm} (\text{of-real } x - u)$  by (rule complex-Re-le-cmod)
    also have ...  $= x / 2$  using dist by (simp add: norm-minus-commute)
    finally have  $Re u: Re u \geq x/2$  using elim by simp
    have  $\text{norm} (\text{stirling-integral } m u / (u - \text{of-real } x) \wedge \text{Suc } j) \leq$ 
         $c / Re u \wedge m / (x / 2) \wedge \text{Suc } j$  using Re-u elim
      unfolding norm-divide norm-power dist
      by (intro divide-right-mono zero-le-power c) simp-all
    also have ...  $\leq c / (x/2) \wedge m / (x / 2) \wedge \text{Suc } j$  using c-nonneg elim Re-u
      by (intro divide-right-mono divide-left-mono power-mono) simp-all
    also have ...  $= B / x \wedge (m + \text{Suc } j)$  using elim by (simp add: B-def field-simps power-add)
    finally show  $\text{norm} (\text{stirling-integral } m u / (u - \text{of-real } x) \wedge \text{Suc } j) \leq B / x$ 
       $\wedge (m + \text{Suc } j)$ .
  qed (insert elim c-nonneg, auto simp: B-def simp del: power-Suc)
  hence  $\text{cmod} ((\text{deriv} \wedge j) (\text{stirling-integral } m) (\text{of-real } x)) \leq B' / x \wedge (j + m)$ 
    using elim by (simp add: field-simps norm-divide norm-mult norm-power B'-def)
  with elim m show ?case by (simp-all add: add-ac deriv-stirling-integral-complex-of-real)
  qed
  thus ?thesis by (rule bigoI)
  qed

definition stirling-sum where
  stirling-sum  $j m x =$ 
     $(-1) \wedge j * (\sum k = 1..<m. (\text{of-real } (\text{beroulli } (\text{Suc } k)) * \text{pochhammer } (\text{of-nat } k) j / (\text{of-nat } k * \text{of-nat } (\text{Suc } k))) * \text{inverse } x \wedge (k + j))$ 

definition stirling-sum' where
  stirling-sum'  $j m x =$ 

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$$(-1) \wedge (\text{Suc } j) * (\sum k \leq m. (\text{of-real} (\text{bernoulli}' k) * \text{pochhammer} (\text{of-nat} (\text{Suc } k)) (j - 1) * \text{inverse } x \wedge (k + j)))$$

**lemma** *stirling-sum-complex-of-real*:

*stirling-sum j m (complex-of-real x) = complex-of-real (stirling-sum j m x)*  
**by** (*simp add: stirling-sum-def pochhammer-of-real [symmetric] del: of-nat-Suc*)

**lemma** *stirling-sum'-complex-of-real*:

*stirling-sum' j m (complex-of-real x) = complex-of-real (stirling-sum' j m x)*  
**by** (*simp add: stirling-sum'-def pochhammer-of-real [symmetric] del: of-nat-Suc*)

**lemma** *has-field-derivative-stirling-sum-complex [derivative-intros]*:

*Re x > 0  $\implies$  (stirling-sum j m has-field-derivative stirling-sum (Suc j) m x) (at x)*

**unfolding** *stirling-sum-def [abs-def] sum-distrib-left*

**by** (*rule DERIV-sum*) (*auto intro!: derivative-eq-intros simp del: of-nat-Suc simp: pochhammer-Suc power-diff*)

**lemma** *has-field-derivative-stirling-sum-real [derivative-intros]*:

*x > (0::real)  $\implies$  (stirling-sum j m has-field-derivative stirling-sum (Suc j) m x) (at x)*

**unfolding** *stirling-sum-def [abs-def] sum-distrib-left*

**by** (*rule DERIV-sum*) (*auto intro!: derivative-eq-intros simp del: of-nat-Suc simp: pochhammer-Suc power-diff*)

**lemma** *has-field-derivative-stirling-sum'-complex [derivative-intros]*:

**assumes** *j > 0 Re x > 0*

**shows** *(stirling-sum' j m has-field-derivative stirling-sum' (Suc j) m x) (at x)*

**proof** (*cases j*)

**case** (*Suc j'*)

**from** *assms have [simp]: x ≠ 0* **by** *auto*

**define** *c where c = (λn. (-1) ^ Suc j \* complex-of-real (bernoulli' n) \* pochhammer (of-nat (Suc n)) j')*

**define** *T where T = (λn x. c n \* inverse x ^ (j + n))*

**define** *T' where T' = (λn x. - (of-nat (j + n)) \* c n \* inverse x ^ (Suc (j + n)))*

**have** *((λx. ∑ k ≤ m. T k x) has-field-derivative (∑ k ≤ m. T' k x)) (at x)* **using** *assms Suc*

**by** (*intro DERIV-sum*)

*(auto simp: T-def T'-def intro!: derivative-eq-intros*

*simp: field-simps power-add [symmetric] simp del: of-nat-Suc power-Suc of-nat-add)*

**also have** *(λx. (∑ k ≤ m. T k x)) = stirling-sum' j m*

**by** (*simp add: Suc T-def c-def stirling-sum'-def fun-eq-iff add-ac mult.assoc sum-distrib-left*)

**also have** *(∑ k ≤ m. T' k x) = stirling-sum' (Suc j) m x*

**by** (*simp add: T'-def c-def Suc stirling-sum'-def sum-distrib-left sum-distrib-right algebra-simps pochhammer-Suc*)

**finally show** *?thesis* .

```

qed (insert assms, simp-all)

lemma has-field-derivative-stirling-sum'-real [derivative-intros]:
  assumes j > 0 x > (0::real)
  shows (stirling-sum' j m has-field-derivative stirling-sum' (Suc j) m x) (at x)
proof (cases j)
  case (Suc j')
  from assms have [simp]: x ≠ 0 by auto
  define c where c = (λn. (-1) ^ Suc j * (bernoulli' n) * pochhammer (of-nat (Suc n)) j')
  define T where T = (λn x. c n * inverse x ^ (j + n))
  define T' where T' = (λn x. - (of-nat (j + n)) * c n * inverse x ^ (Suc (j + n)))
  have ((λx. ∑ k≤m. T k x) has-field-derivative (∑ k≤m. T' k x)) (at x) using
    assms Suc
    by (intro DERIV-sum)
    (auto simp: T-def T'-def intro!: derivative-eq-intros
      simp: field-simps power-add [symmetric] simp del: of-nat-Suc power-Suc
      of-nat-add)
  also have (λx. (∑ k≤m. T k x)) = stirling-sum' j m
    by (simp add: Suc T-def c-def stirling-sum'-def fun-eq-iff add-ac mult.assoc
      sum-distrib-left)
  also have (∑ k≤m. T' k x) = stirling-sum' (Suc j) m x
    by (simp add: T'-def c-def Suc stirling-sum'-def sum-distrib-left
      sum-distrib-right algebra-simps pochhammer-Suc)
  finally show ?thesis .
qed (insert assms, simp-all)

lemma higher-deriv-stirling-sum-complex:
  Re x > 0 ==> (deriv ^ i) (stirling-sum j m) x = stirling-sum (i + j) m x
proof (induction i arbitrary: x)
  case (Suc i)
  have deriv ((deriv ^ i) (stirling-sum j m)) x = deriv (stirling-sum (i + j) m) x
    using eventually-nhds-in-open[of {x. Re x > 0} x] Suc.prems
    by (intro deriv-cong-ev refl) (auto elim!: eventually-mono simp: open-halfspace-Re-gt
      Suc.IH)
  also from Suc.prems have ... = stirling-sum (Suc (i + j)) m x
    by (intro DERIV-imp-deriv has-field-derivative-stirling-sum-complex)
  finally show ?case by simp
qed simp-all

definition Polygamma-approx :: nat ⇒ nat ⇒ 'a ⇒ 'a :: {real-normed-field, ln}
where
  Polygamma-approx j m =
    (deriv ^ j) (λx. (x - 1 / 2) * ln x - x + of-real (ln (2 * pi)) / 2 +
    stirling-sum 0 m x)

lemma Polygamma-approx-Suc: Polygamma-approx (Suc j) m = deriv (Polygamma-approx

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 $j m)$ 
by (simp add: Polygamma-approx-def)

lemma Polygamma-approx-0:
  Polygamma-approx 0 m x = (x - 1/2) * ln x - x + of-real (ln (2*pi)) / 2 + stirling-sum 0 m x
by (simp add: Polygamma-approx-def)

lemma Polygamma-approx-1-complex:
  Re x > 0 ==>
  Polygamma-approx (Suc 0) m x = ln x - 1 / (2*x) + stirling-sum (Suc 0) m x
unfoldng Polygamma-approx-Suc Polygamma-approx-0
by (intro DERIV-imp-deriv)
  (auto intro!: derivative-eq-intros elim!: nonpos-Reals-cases simp: field-simps)

lemma Polygamma-approx-1-real:
  x > (0 :: real) ==>
  Polygamma-approx (Suc 0) m x = ln x - 1 / (2*x) + stirling-sum (Suc 0) m x
unfoldng Polygamma-approx-Suc Polygamma-approx-0
by (intro DERIV-imp-deriv)
  (auto intro!: derivative-eq-intros elim!: nonpos-Reals-cases simp: field-simps)

lemma stirling-sum-2-conv-stirling-sum'-1:
  fixes x :: 'a :: {real-div-algebra, field-char-0}
  assumes m > 0 x ≠ 0
  shows stirling-sum' 1 m x = 1 / x + 1 / (2 * x^2) + stirling-sum 2 m x
proof -
  have pochhammer-2: pochhammer (of-nat k) 2 = of-nat k * of-nat (Suc k) for k
    by (simp add: pochhammer-Suc eval-nat-numeral add-ac)
  have stirling-sum 2 m x =
     $(\sum k = Suc 0 .. < m. of-real (bernoulli' (Suc k)) * inverse x \wedge Suc (Suc k))$ 
  unfoldng stirling-sum-def pochhammer-2 power2-minus power-one mult-1-left
  by (intro sum.cong refl)
    (simp-all add: stirling-sum-def pochhammer-2 power2-eq-square divide-simps bernoulli'-def
      del: of-nat-Suc power-Suc)
  also have  $1 / (2 * x^2) + \dots =$ 
     $(\sum k=0 .. < m. of-real (bernoulli' (Suc k)) * inverse x \wedge Suc (Suc k))$ 
  using assms
  by (subst (2) sum.atLeast-Suc-lessThan) (simp-all add: power2-eq-square field-simps)
  also have  $1 / x + \dots = (\sum k=0 .. < Suc m. of-real (bernoulli' k) * inverse x \wedge Suc k)$ 
    by (subst sum.atLeast0-lessThan-Suc-shift) (simp-all add: bernoulli'-def divide-simps)
  also have  $\dots = (\sum k \leq m. of-real (bernoulli' k) * inverse x \wedge Suc k)$ 
    by (intro sum.cong auto)

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also have ... = stirling-sum' 1 m x by (simp add: stirling-sum'-def)
finally show ?thesis by (simp add: add-ac)
qed

lemma Polygamma-approx-2-real:
assumes x > (0::real) m > 0
shows Polygamma-approx (Suc (Suc 0)) m x = stirling-sum' 1 m x
proof -
have Polygamma-approx (Suc (Suc 0)) m x = deriv (Polygamma-approx (Suc
0) m) x
by (simp add: Polygamma-approx-Suc)
also have ... = deriv (λx. ln x - 1 / (2*x) + stirling-sum (Suc 0) m x) x
using eventually-nhds-in-open[of {0<..} x] assms
by (intro deriv-cong-ev) (auto elim!: eventually-mono simp: Polygamma-approx-1-real)
also have ... = 1 / x + 1 / (2*x^2) + stirling-sum (Suc (Suc 0)) m x using
assms
by (intro DERIV-imp-deriv) (auto intro!: derivative-eq-intros
elim!: nonpos-Reals-cases simp: field-simps power2-eq-square)
also have ... = stirling-sum' 1 m x using stirling-sum-2-conv-stirling-sum'-1[of
m x] assms
by (simp add: eval-nat-numeral)
finally show ?thesis .
qed

lemma Polygamma-approx-2-complex:
assumes Re x > 0 m > 0
shows Polygamma-approx (Suc (Suc 0)) m x = stirling-sum' 1 m x
proof -
have Polygamma-approx (Suc (Suc 0)) m x = deriv (Polygamma-approx (Suc
0) m) x
by (simp add: Polygamma-approx-Suc)
also have ... = deriv (λx. ln x - 1 / (2*x) + stirling-sum (Suc 0) m x) x
using eventually-nhds-in-open[of {s. Re s > 0} x] assms
by (intro deriv-cong-ev)
(auto simp: open-halfspace-Re-gt elim!: eventually-mono simp: Polygamma-approx-1-complex)
also have ... = 1 / x + 1 / (2*x^2) + stirling-sum (Suc (Suc 0)) m x using
assms
by (intro DERIV-imp-deriv) (auto intro!: derivative-eq-intros
elim!: nonpos-Reals-cases simp: field-simps power2-eq-square)
also have ... = stirling-sum' 1 m x using stirling-sum-2-conv-stirling-sum'-1[of
m x] assms
by (subst stirling-sum-2-conv-stirling-sum'-1) (auto simp: eval-nat-numeral)
finally show ?thesis .
qed

lemma Polygamma-approx-ge-2-real:
assumes x > (0::real) m > 0
shows Polygamma-approx (Suc (Suc j)) m x = stirling-sum' (Suc j) m x
using assms(1)

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proof (induction j arbitrary: x)
  case ( $0\ x$ )
    with assms show ?case by (simp add: Polygamma-approx-2-real)
  next
    case ( $\text{Suc } j\ x$ )
      have Polygamma-approx ( $\text{Suc } (\text{Suc } (\text{Suc } j))\ m\ x = \text{deriv } (\text{Polygamma-approx}$ 
      ( $\text{Suc } (\text{Suc } j))\ m\ x$ )
        by (simp add: Polygamma-approx-Suc)
      also have ... = deriv (stirling-sum' ( $\text{Suc } j$ )  $m\ x$ )
        using eventually-nhds-in-open[of { $0 < ..$ }  $x$ ] Suc.prems
        by (intro deriv-cong-ev refl) (auto elim!: eventually-mono simp: Suc.IH)
      also have ... = stirling-sum' ( $\text{Suc } (\text{Suc } j)$   $m\ x$ ) using Suc.prems
        by (intro DERIV-imp-deriv derivative-intros) simp-all
    finally show ?case .
  qed

lemma Polygamma-approx-ge-2-complex:
  assumes Re x > 0 m > 0
  shows Polygamma-approx ( $\text{Suc } (\text{Suc } j)\ m\ x = \text{stirling-sum}' (\text{Suc } j)\ m\ x$ )
  using assms(1)
  proof (induction j arbitrary: x)
    case ( $0\ x$ )
      with assms show ?case by (simp add: Polygamma-approx-2-complex)
  next
    case ( $\text{Suc } j\ x$ )
      have Polygamma-approx ( $\text{Suc } (\text{Suc } (\text{Suc } j))\ m\ x = \text{deriv } (\text{Polygamma-approx}$ 
      ( $\text{Suc } (\text{Suc } j))\ m\ x$ )
        by (simp add: Polygamma-approx-Suc)
      also have ... = deriv (stirling-sum' ( $\text{Suc } j$ )  $m\ x$ )
        using eventually-nhds-in-open[of { $x. \text{Re } x > 0$ }  $x$ ] Suc.prems
        by (intro deriv-cong-ev refl) (auto elim!: eventually-mono simp: Suc.IH open-halfspace-Re-gt)
      also have ... = stirling-sum' ( $\text{Suc } (\text{Suc } j)\ m\ x$ ) using Suc.prems
        by (intro DERIV-imp-deriv derivative-intros) simp-all
    finally show ?case .
  qed

lemma Polygamma-approx-complex-of-real:
  assumes  $x > 0\ m > 0$ 
  shows Polygamma-approx  $j\ m$  (complex-of-real  $x$ ) = of-real (Polygamma-approx
   $j\ m\ x$ )
  proof (cases j)
    case  $0$ 
      with assms show ?thesis by (simp add: Polygamma-approx-0 Ln-of-real stir-
      ling-sum-complex-of-real)
  next
    case [simp]: ( $\text{Suc } j'$ )
      thus ?thesis
    proof (cases j')
      case  $0$ 

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with assms show ?thesis
  by (simp add: Polygamma-approx-1-complex
    Polygamma-approx-1-real stirling-sum-complex-of-real Ln-of-real)
next
  case (Suc j")
  with assms show ?thesis
    by (simp add: Polygamma-approx-ge-2-complex Polygamma-approx-ge-2-real
      stirling-sum'-complex-of-real)
qed
qed

lemma higher-deriv-Polygamma-approx [simp]:
  (deriv ^~ j) (Polygamma-approx i m) = Polygamma-approx (j + i) m
  by (simp add: Polygamma-approx-def funpow-add)

lemma stirling-sum-holomorphic [holomorphic-intros]:
  0 ∉ A ⟹ stirling-sum j m holomorphic-on A
  unfolding stirling-sum-def by (intro holomorphic-intros) auto

lemma Polygamma-approx-holomorphic [holomorphic-intros]:
  Polygamma-approx j m holomorphic-on {s. Re s > 0}
  unfolding Polygamma-approx-def
  by (intro holomorphic-intros) (auto simp: open-halfspace-Re-gt elim!: nonpos-Reals-cases)

lemma higher-deriv-lnGamma-stirling:
  assumes m: m > 0
  shows (λx::real. (deriv ^~ j) ln-Gamma x - Polygamma-approx j m x) ∈ O(λx.
  1 / x ^ (m + j))
proof -
  have eventually (λx. |(deriv ^~ j) ln-Gamma x - Polygamma-approx j m x| =
    inverse (real m) * |(deriv ^~ j) (stirling-integral m) x|) at-top
  using eventually-gt-at-top[of 0::real]
proof eventually-elim
  case (elim x)
  note x = this
  have ∀ F y in nhds (complex-of-real x). y ∈ - ℝ≤0
    using elim by (intro eventually-nhds-in-open) auto
  hence (deriv ^~ j) (λx. ln-Gamma x - Polygamma-approx 0 m x) (complex-of-real
  x) =
    (deriv ^~ j) (λx. (-inverse (of-nat m)) * stirling-integral m x)
  (complex-of-real x)
  using x m
  by (intro higher-deriv-cong-ev refl)
  (auto elim!: eventually-mono simp: ln-Gamma-stirling-complex Polygamma-approx-def
    field-simps open-halfspace-Re-gt stirling-sum-def)
  also have ... = - inverse (of-nat m) * (deriv ^~ j) (stirling-integral m) (of-real
  x)
  using x m
  by (intro higher-deriv-cmult[of - -ℝ≤0] stirling-integral-holomorphic)

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(auto simp: open-halfspace-Re-gt)
also have (deriv  $\wedge\wedge j$ ) ( $\lambda x. \ln\text{-}\Gamma x - \text{Polygamma-approx } 0 m x$ ) (complex-of-real
 $x) = 
(deriv  $\wedge\wedge j$ ) \ln\text{-}\Gamma (\text{of-real } x) - (deriv  $\wedge\wedge j$ ) (\text{Polygamma-approx }
0 m) (\text{of-real } x)
using x
by (intro higher-deriv-diff[of - {s. Re s > 0}])
(auto intro!: holomorphic-intros elim!: nonpos-Real-cases simp: open-halfspace-Re-gt)
also have (deriv  $\wedge\wedge j$ ) (\text{Polygamma-approx } 0 m) (complex-of-real x) =
of-real (\text{Polygamma-approx } j m x) using x m
by (simp add: Polygamma-approx-complex-of-real)
also have norm (- inverse (of-nat m) * (deriv  $\wedge\wedge j$ ) (\text{stirling-integral } m)
(complex-of-real x)) =
inverse (real m) * |(deriv  $\wedge\wedge j$ ) (\text{stirling-integral } m) x|
using x m by (simp add: norm-mult norm-inverse deriv-stirling-integral-complex-of-real)
also have (deriv  $\wedge\wedge j$ ) \ln\text{-}\Gamma (complex-of-real x) = of-real ((deriv  $\wedge\wedge j$ )
\ln\text{-}\Gamma x) using x
by (simp add: higher-deriv-\ln\text{-}\Gamma-complex-of-real)
also have norm (... - of-real (\text{Polygamma-approx } j m x)) =
|(deriv  $\wedge\wedge j$ ) \ln\text{-}\Gamma x - \text{Polygamma-approx } j m x|
by (simp only: of-real-diff [symmetric] norm-of-real)
finally show ?case .
qed
from bithetaI-cong[OF this] m
have ( $\lambda x::\text{real}. (deriv \wedge\wedge j) \ln\text{-}\Gamma x - \text{Polygamma-approx } j m x \in
\Theta(\lambda x. (deriv \wedge\wedge j) (\text{stirling-integral } m) x)$ ) by simp
also have ( $\lambda x::\text{real}. (deriv \wedge\wedge j) (\text{stirling-integral } m) x \in O(\lambda x. 1 / x^{\wedge(m + j)})$ ) using m
by (rule deriv-stirling-integral-real-bound)
finally show ?thesis .
qed

lemma Polygamma-approx-1-real':
assumes x:  $(x::\text{real}) > 0$  and m:  $m > 0$ 
shows  $\text{Polygamma-approx } 1 m x = \ln x - (\sum k = \text{Suc } 0..m. \text{beroulli}' k * \text{inverse } x^{\wedge k} / \text{real } k)$ 
proof -
have Polygamma-approx 1 m x =  $\ln x - (1 / (2 * x) + (\sum k=\text{Suc } 0..<m. \text{beroulli}' (\text{Suc } k) * \text{inverse } x^{\wedge \text{Suc } k} / \text{real } (\text{Suc } k)))$ 
(is - = - - (- + ?S)) using x by (simp add: Polygamma-approx-1-real stirling-sum-def)
also have ?S =  $(\sum k=\text{Suc } 0..<m. \text{beroulli}' (\text{Suc } k) * \text{inverse } x^{\wedge \text{Suc } k} / \text{real } (\text{Suc } k))$ 
by (intro sum.cong refl) (simp-all add: beroulli'-def)
also have  $1 / (2 * x) + \dots = (\sum k=0..<m. \text{beroulli}' (\text{Suc } k) * \text{inverse } x^{\wedge \text{Suc } k} / \text{real } (\text{Suc } k))$ 
using m
by (subst (2) sum.atLeast-Suc-lessThan) (simp-all add: field-simps)
also have ... =  $(\sum k = \text{Suc } 0..m. \text{beroulli}' k * \text{inverse } x^{\wedge k} / \text{real } k)$  using$ 
```

*assms*  
**by** (*subst sum.shift-bounds-Suc-ivl [symmetric]*) (*simp add: atLeastLessThanSuc-atLeastAtMost*)  
**finally show** ?thesis .  
**qed**

**theorem**  
**assumes**  $m: m > 0$   
**shows** *ln-Gamma-real-asymptotics*:  

$$(\lambda x. \ln\text{-}\Gamma x - ((x - 1 / 2) * \ln x - x + \ln(2 * \pi) / 2 + (\sum_{k=1..m} \text{bernoulli}(Suc k) / (\text{real } k * \text{real } (\text{Suc } k)) / x^k))) \in O(\lambda x. 1 / x^m)$$
 (**is** ?th1)

**and** *Digamma-real-asymptotics*:  

$$(\lambda x. \text{Digamma } x - (\ln x - (\sum_{k=1..m} \text{bernoulli}' k / \text{real } k / x^k))) \in O(\lambda x. 1 / (x^{\text{Suc } m}))$$
 (**is** ?th2)

**and** *Polygamma-real-asymptotics*:  $j > 0 \implies$   

$$(\lambda x. \text{Polygamma } j x - (-1)^j \sum_{k \leq m} \text{bernoulli}' k * \text{pochhammer}(\text{real } (\text{Suc } k), (j - 1) / x^{k+j})) \in O(\lambda x. 1 / x^{(m+j+1)})$$
 (**is** -  $\implies$  ?th3)

**proof** –  
**define**  $G :: nat \Rightarrow real \Rightarrow real$  **where**  
 $G = (\lambda m. \text{if } m = 0 \text{ then } \ln\text{-}\Gamma \text{ else } \text{Polygamma } (m - 1))$   
**have**  $*: (\lambda x. G j x - h x) \in O(\lambda x. 1 / x^{(m+j)})$   
**if**  $\bigwedge_{x:\text{real}} x > 0 \implies \text{Polygamma-approx } j m x = h x$  **for**  $j h$   
**proof** –  
**have**  $(\lambda x. G j x - h x) \in \Theta(\lambda x. (\text{deriv } \wedge j) \ln\text{-}\Gamma x - \text{Polygamma-approx } j m x)$  (**is** -  $\in \Theta(?f)$ )  
**using** that  
**by** (*intro bigthetaI-cong*) (*auto intro: eventually-mono[*OF eventually-gt-at-top[*of 0::real*]]*)  
*simp del: funpow.simps simp: higher-deriv-ln-Gamma-real G-def*  
**also have**  $?f \in O(\lambda x:\text{real}. 1 / x^{(m+j)})$  **using**  $m$   
**by** (*rule higher-deriv-lnGamma-stirling*)  
**finally show** ?thesis .  
**qed***

**note** [[*simproc del: simplify-landau-sum*]]  
**from** \*[*OF Polygamma-approx-0*] *assms* **show** ?th1  
**by** (*simp add: G-def Polygamma-approx-0 stirling-sum-def field-simps*)  
**from** \*[*OF Polygamma-approx-1-real*] *assms* **show** ?th2 **by** (*simp add: G-def field-simps*)

**assume**  $j: j > 0$   
**from** \*[*OF Polygamma-approx-ge-2-real, of j - 1*] *assms j* **show** ?th3  
**by** (*simp add: G-def stirling-sum'-def power-add power-diff field-simps*)  
**qed**

## 2.5 Asymptotics of the complex Gamma function

The  $m$ -th order remainder of Stirling's formula for  $\log \Gamma$  is  $O(s^{-m})$  uniformly over any complex cone  $\text{Arg}(z) \leq \alpha$ ,  $z \neq 0$  for any angle  $\alpha \in (0, \pi)$ . This means that there is bounded by  $cz^{-m}$  for some constant  $c$  for all  $z$  in this cone.

**context**

fixes  $F$  and  $\alpha$

assumes  $\alpha: \alpha \in \{0 <.. < pi\}$

defines  $F \equiv \text{principal}(\text{complex-cone}' \alpha - \{0\})$

begin

**lemma** *stirling-integral-bigo*:

fixes  $m :: \text{nat}$

assumes  $m: m > 0$

shows *stirling-integral*  $m \in O[F](\lambda s. 1 / s^m)$

**proof** –

obtain  $c$  where  $c: \bigwedge s. s \in \text{complex-cone}' \alpha - \{0\} \implies \text{norm}(\text{stirling-integral } m s) \leq c / \text{norm } s^m$

using *stirling-integral-bound'*[ $OF \langle m > 0 \rangle \alpha$ ] by *blast*

have  $0 \leq \text{norm}(\text{stirling-integral } m 1 :: \text{complex})$

by *simp*

also have ...  $\leq c$

using  $c[\text{of } 1] \alpha$  by *simp*

finally have  $c \geq 0$ .

have *eventually*  $(\lambda s. s \in \text{complex-cone}' \alpha - \{0\}) F$

unfolding  $F\text{-def}$  by (*auto simp: eventually-principal*)

hence *eventually*  $(\lambda s. \text{norm}(\text{stirling-integral } m s) \leq c * \text{norm}(1 / s^m)) F$

by *eventually-elim* (*use c in <simp add: norm-divide norm-power>*)

thus *stirling-integral*  $m \in O[F](\lambda s. 1 / s^m)$

by (*intro bigoI[of - c]*) *auto*

**qed**

**end**

The following is a more explicit statement of this:

**theorem** *ln-Gamma-complex-asymptotics-explicit*:

fixes  $m :: \text{nat}$  and  $\alpha :: \text{real}$

assumes  $m > 0$  and  $\alpha \in \{0 <.. < pi\}$

obtains  $C :: \text{real}$  and  $R :: \text{complex} \Rightarrow \text{complex}$

where  $\forall s :: \text{complex}. s \notin \mathbb{R}_{\leq 0} \rightarrow$

$\text{ln-Gamma } s = (s - 1/2) * \ln s - s + \ln(2 * pi) / 2 +$

$(\sum_{k=1..<m. \text{bernonulli}} (k+1) / (k * (k+1) * s^k)) - R s$

and  $\forall s. s \neq 0 \wedge |\text{Arg } s| \leq \alpha \rightarrow \text{norm}(R s) \leq C / \text{norm } s^m$

**proof** –

obtain  $c$  where  $c: \bigwedge s. s \in \text{complex-cone}' \alpha - \{0\} \implies \text{norm}(\text{stirling-integral } m s) \leq c / \text{norm } s^m$

```

using stirling-integral-bound'[OF assms] by blast
have 0 ≤ norm (stirling-integral m 1 :: complex)
    by simp
also have ... ≤ c
    using c[of 1] assms by simp
finally have c ≥ 0 .
define R where R = (λs::complex. stirling-integral m s / of-nat m)
show ?thesis
proof (rule that)
    from ln-Gamma-stirling-complex[of - m] assms show
        ∀ s::complex. s ∉ ℝ≤₀ →
            ln-Gamma s = (s - 1 / 2) * ln s - s + ln (2 * pi) / 2 +
            (∑ k=1... bernoulli (k+1) / (k * (k+1) * s ^ k)) - R s
        by (auto simp add: R-def algebra-simps)
    show ∀ s. s ≠ 0 ∧ |Arg s| ≤ α → cmod (R s) ≤ c / real m / cmod s ^ m
    proof (safe, goal-cases)
        case (1 s)
        show ?case
            using 1 c[of s] assms
            by (auto simp: complex-cone-altdef abs-le-iff R-def norm-divide field-simps)
    qed
    qed
qed

```

Lastly, we can also derive the asymptotics of  $\Gamma$  itself:

$$\Gamma(z) \sim \sqrt{2\pi/z} \left(\frac{z}{e}\right)^z$$

uniformly for  $|z| \rightarrow \infty$  within the cone  $\text{Arg}(z) \leq \alpha$  for  $\alpha \in (0, \pi)$ :

```

context
fixes F and α
assumes α: α ∈ {0 <.. < pi}
defines F ≡ inf at-infinity (principal (complex-cone' α))
begin

lemma Gamma-complex-asymp-equiv:
    Gamma ~[F] (λs. sqrt (2 * pi) * (s / exp 1) powr s / s powr (1 / 2))
proof –
    define I :: complex ⇒ complex where I = stirling-integral 1
    have eventually (λs. s ∈ complex-cone' α) F
        by (auto simp: eventually-inf-principal F-def)
    moreover have eventually (λs. s ≠ 0) F
        unfold F-def eventually-inf-principal
        using eventually-not-equal-at-infinity by eventually-elim auto
    ultimately have eventually (λs. Gamma s =
        sqrt (2 * pi) * (s / exp 1) powr s / s powr (1 / 2) / exp (I s)) F
proof eventually-elim
    case (elim s)
    from elim have s': s ∉ ℝ≤₀

```

```

using complex-cone-inter-nonpos-Reals[of  $-\alpha \alpha$ ]  $\alpha$  by auto
from elim have [simp]:  $s \neq 0$  by auto
from  $s'$  have  $\text{Gamma } s = \exp(\ln\text{-}\Gamma s)$ 
  unfolding  $\text{Gamma-complex-altdef}$  using nonpos-Ints-subset-nonpos-Reals by
auto
  also from  $s'$  have  $\ln\text{-}\Gamma s = (s - 1/2) * \ln s - s + \text{complex-of-real}(\ln(2 * \pi) / 2) - I s$ 
    by (subst ln-Gamma-stirling-complex[of - 1]) (simp-all add: exp-add exp-diff
I-def)
  also have  $\exp \dots = \exp((s - 1/2) * \ln s) / \exp s *$ 
     $\exp(\text{complex-of-real}(\ln(2 * \pi) / 2)) / \exp(I s)$ 
    unfolding exp-diff exp-add by (simp add: exp-diff exp-add)
  also have  $\exp((s - 1/2) * \ln s) = s \text{ powr}(s - 1/2)$ 
    by (simp add: powr-def)
  also have  $\exp(\text{complex-of-real}(\ln(2 * \pi) / 2)) = \sqrt(2 * \pi)$ 
    by (subst exp-of-real) (auto simp: powr-def simp flip: powr-half-sqrt)
  also have  $\exp s = \exp 1 \text{ powr } s$ 
    by (simp add: powr-def)
  also have  $s \text{ powr}(s - 1/2) / \exp 1 \text{ powr } s = (s \text{ powr } s / \exp 1 \text{ powr } s) / s$ 
    powr(1/2)
    by (subst powr-diff) auto
  also have  $*: \ln(s / \exp 1) = \ln s - 1$ 
    using Ln-divide-of-real[of exp 1 s] by (simp flip: exp-of-real)
  hence  $s \text{ powr } s / \exp 1 \text{ powr } s = (s / \exp 1) \text{ powr } s$ 
    unfolding powr-def by (subst *) (auto simp: exp-diff field-simps)
  finally show  $\text{Gamma } s = \sqrt(2 * \pi) * (s / \exp 1) \text{ powr } s / s \text{ powr } (1 / 2)$ 
  /  $\exp(I s)$ 
    by (simp add: algebra-simps)
qed
hence  $\text{Gamma} \sim[F] (\lambda s. \sqrt(2 * \pi) * (s / \exp 1) \text{ powr } s / s \text{ powr } (1 / 2) /$ 
 $\exp(I s))$ 
  by (rule asymp-equiv-refl-ev)
also have  $\dots \sim[F] (\lambda s. \sqrt(2 * \pi) * (s / \exp 1) \text{ powr } s / s \text{ powr } (1 / 2) / 1)$ 
proof (intro asymp-equiv-intros)
  have  $F \leq \text{principal}(\text{complex-cone}' \alpha - \{0\})$ 
    unfolding le-principal F-def eventually-inf-principal
    using eventually-not-equal-at-infinity by eventually-elim auto
  moreover have  $I \in O[\text{principal}(\text{complex-cone}' \alpha - \{0\})](\lambda s. 1 / s)$ 
    using stirling-integral-bigO[of  $\alpha 1$ ]  $\alpha$  unfolding F-def by (simp add: I-def)
  ultimately have  $I \in O[F](\lambda s. 1 / s)$ 
    by (rule landau-o.big.filter-mono)
  also have  $(\lambda s. 1 / s) \in o[F](\lambda s. 1)$ 
proof (rule landau-o.smallI)
  fix  $c :: \text{real}$ 
  assume  $c: c > 0$ 
  hence eventually  $(\lambda z::\text{complex}. \text{norm } z \geq 1 / c)$  at-infinity
    by (auto simp: eventually-at-infinity)
  moreover have eventually  $(\lambda z::\text{complex}. z \neq 0)$  at-infinity
    by (rule eventually-not-equal-at-infinity)

```

```

ultimately show eventually ( $\lambda z::\text{complex}. \text{norm} (1 / z) \leq c * \text{norm} (1 :: \text{complex})$ )  $F$ 
  unfolding  $F$ -def eventually-inf-principal
  by eventually-elim (use  $\langle c > 0 \rangle$  in  $\langle \text{auto simp: norm-divide field-simps} \rangle$ )
  qed
  finally have  $I \in o[F](\lambda s. 1)$  .
  from smalloD-tendsto[OF this] have [tendsto-intros]:  $(I \longrightarrow 0) F$ 
    by simp
    show  $(\lambda x. \exp(I x)) \sim[F] (\lambda x. 1)$ 
      by (rule asymp-equivI' tendsto-eq-intros refl | simp)+
    qed
    finally show ?thesis by simp
  qed

end

end

```

## References

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- [3] G. J. O. Jameson. A simple proof of Stirling's formula for the Gamma function. *The Mathematical Gazette*, 99:68–74, 3 2015.