

Formalizing Statecharts using Hierarchical Automata

Steffen Helke and Florian Kammüller

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Abstract

We formalize in Isabelle/HOL the abstract syntax and a synchronous step semantics for the specification language Statecharts [HN96]. The formalization is based on Hierarchical Automata [MLS97] which allow a structural decomposition of Statecharts into Sequential Automata. To support the composition of Statecharts, we introduce calculating operators to construct a Hierarchical Automaton in a stepwise manner [HK01]. Furthermore, we present a complete semantics of Statecharts including a theory of data spaces, which enables the modelling of racing effects [HK05]. We also adapt CTL for Statecharts to build a bridge for future combinations with model checking. However the main motivation of this work is to provide a sound and complete basis for reasoning on Statecharts. As a central meta theorem we prove that the well-formedness of a Statechart is preserved by the semantics [Hel07].

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1 Contributions to the Standard Library of HOL

```
theory Contrib
imports Main
begin
```

1.1 Basic definitions and lemmas

1.1.1 Lambda expressions

```
definition restrict :: ['a => 'b, 'a set] => ('a => 'b) where
  restrict f A = (%x. if x : A then f x else (@ y. True))
```

```
syntax (ASCII)
```

```
-lambda-in :: [pttrn, 'a set, 'a => 'b] => ('a => 'b) ((%%-:./ -) [0, 0, 3] 3)
```

```
syntax
```

```
-lambda-in :: [pttrn, 'a set, 'a => 'b] => ('a => 'b) ((%λ-∈-./ -) [0, 0, 3] 3)
```

```
translations
```

```
λx∈A. f == CONST restrict (%x. f) A
```

1.1.2 Maps

```
definition chg-map :: ('b => 'b) => 'a => ('a → 'b) => ('a → 'b) where
  chg-map f a m = (case m a of None => m | Some b => m(a|→f b))
```

```
lemma map-some-list [simp]:
```

```
  map the (map Some L) = L
```

```
apply (induct-tac L)
```

```
apply auto
```

```
done
```

```
lemma dom-ran-the:
```

```
  [ ran G = {y}; x ∈ (dom G) ] ==> (the (G x)) = y
```

```
apply (unfold ran-def dom-def)
```

```
apply auto
```

```
done
```

```
lemma dom-None:
```

$(S \notin \text{dom } F) \implies (F S = \text{None})$
by (*unfold dom-def, auto*)

lemma *ran-dom-the*:

$\llbracket y \notin \text{Union } (\text{ran } G); x \in \text{dom } G \rrbracket \implies y \notin \text{the } (G x)$
by (*unfold ran-def dom-def, auto*)

lemma *dom-map-upd*: $\text{dom}(m(a|-\>b)) = \text{insert } a (\text{dom } m)$
apply *auto*
done

1.1.3 *rtrancl*

lemma *rtrancl-Int*:

$\llbracket (a,b) \in A; (a,b) \in B \rrbracket \implies (a,b) \in (A \cap B)^{\hat{*}}$
by *auto*

lemma *rtrancl-mem-Sigma*:

$\llbracket a \neq b; (a,b) \in (A \times A)^{\hat{*}} \rrbracket \implies b \in A$
apply (*frule rtranclD*)
apply (*cut-tac r=A \times A and A=A in trancl-subset-Sigma*)
apply *auto*
done

lemma *help-rtrancl-Range*:

$\llbracket a \neq b; (a,b) \in R^{\hat{*}} \rrbracket \implies b \in \text{Range } R$
apply (*erule rtranclE*)
apply *auto*
done

lemma *rtrancl-Int-help*:

$(a,b) \in (A \cap B)^{\hat{*}} \implies (a,b) \in A^{\hat{*}} \wedge (a,b) \in B^{\hat{*}}$
apply (*unfold Int-def*)
apply *auto*
apply (*rule-tac b=b in rtrancl-induct*)
apply *auto*
apply (*rule-tac b=b in rtrancl-induct*)
apply *auto*
done

lemmas *rtrancl-Int1* = *rtrancl-Int-help* [*THEN conjunct1*]

lemmas *rtrancl-Int2* = *rtrancl-Int-help* [*THEN conjunct2*]

lemma *tranclD3* [*rule-format*]:

$(S,T) \in R^{\hat{+}} \implies (S,T) \notin R \longrightarrow (\exists U. (S,U) \in R \wedge (U,T) \in R^{\hat{+}})$
apply (*rule-tac a=S and b=T and r=R in trancl-induct*)
apply *auto*
done

lemma *tranclD4* [rule-format]:

$(S, T) \in R^+ \implies (S, T) \notin R \longrightarrow (\exists U. (S, U) \in R^+ \wedge (U, T) \in R)$
apply (rule-tac a=S and b=T and r=R in trancl-induct)
apply auto
done

lemma *trancl-collect* [rule-format]:

$\llbracket (x, y) \in R^*; S \notin \text{Domain } R \rrbracket \implies y \neq S \longrightarrow (x, y) \in \{p. \text{fst } p \neq S \wedge \text{snd } p \neq S \wedge p \in R\}^*$
apply (rule-tac b=y in rtrancl-induct)
apply auto
apply (rule rtrancl-into-rtrancl)
apply fast
apply auto
done

lemma *trancl-subseteq*:

$\llbracket R \subseteq Q; S \in R^* \rrbracket \implies S \in Q^*$
apply (frule rtrancl-mono)
apply fast
done

lemma *trancl-Int-subset*:

$(R \cap (A \times A))^+ \subseteq R^+ \cap (A \times A)$
apply (rule subsetI)
apply (rename-tac S)
apply (case-tac S)
apply (rename-tac T V)
apply auto
apply (rule-tac a=T and b=V and r=(R \cap A \times A) in converse-trancl-induct, auto)+
done

lemma *trancl-Int-mem*:

$(S, T) \in (R \cap (A \times A))^+ \implies (S, T) \in R^+ \cap A \times A$
by (rule trancl-Int-subset [THEN subsetD], assumption)

lemma *Int-expand*:

$\{(S, S'). P S S' \wedge Q S S'\} = (\{(S, S'). P S S'\} \cap \{(S, S'). Q S S'\})$
by auto

1.1.4 finite

lemma *finite-conj*:

$\text{finite } (\{(S, S'). P S S'\}::('a*'b)\text{set}) \longrightarrow$
 $\text{finite } \{(S, S'). P (S::'a) (S'::'b) \wedge Q (S::'a) (S'::'b)\}$
apply (rule impI)
apply (subst Int-expand)
apply (rule finite-Int)

apply *auto*
done

lemma *finite-conj2*:
[[*finite A*; *finite B*]] \implies *finite* ({(*S,S'*). *S*: *A* \wedge *S'*: *B*})
by *auto*

1.1.5 *override*

lemma *dom-override-the*:
($x \in (\text{dom } G2)$) \longrightarrow ((*G1* ++ *G2*) *x*) = (*G2* *x*)
by (*auto*)

lemma *dom-override-the2* [*simp*]:
[[*dom G1* \cap *dom G2* = {}]; $x \in (\text{dom } G1)$]] \implies ((*G1* ++ *G2*) *x*) = (*G1* *x*)
apply (*unfold dom-def map-add-def*)
apply *auto*
apply (*drule sym*)
apply (*erule equalityE*)
apply (*unfold Int-def*)
apply *auto*
apply (*erule-tac x=x in allE*)
apply *auto*
done

lemma *dom-override-the3* [*simp*]:
[[$x \notin \text{dom } G2$; $x \in \text{dom } G1$]] \implies ((*G1* ++ *G2*) *x*) = (*G1* *x*)
apply (*unfold dom-def map-add-def*)
apply *auto*
done

lemma *Union-ran-override* [*simp*]:
 $S \in \text{dom } G \implies \bigcup (\text{ran } (G \text{ ++ } \text{Map.empty}(S \mapsto \text{insert } SA (\text{the}(G \ S)))))) =$
 $(\text{insert } SA (\text{Union } (\text{ran } G)))$
apply (*unfold dom-def ran-def*)
apply *auto*
apply (*rename-tac T*)
apply (*case-tac T = S*)
apply *auto*
done

lemma *Union-ran-override2* [*simp*]:
 $S \in \text{dom } G \implies \bigcup (\text{ran } (G(S \mapsto \text{insert } SA (\text{the}(G \ S)))))) = (\text{insert } SA (\text{Union}$
 $(\text{ran } G)))$
apply (*unfold dom-def ran-def*)
apply *auto*
apply (*rename-tac T*)
apply (*case-tac T = S*)
apply *auto*

done

lemma *ran-override* [*simp*]:

$(\text{dom } A \cap \text{dom } B) = \{\} \implies \text{ran } (A ++ B) = (\text{ran } A) \cup (\text{ran } B)$

apply (*unfold Int-def ran-def*)

apply (*simp add: map-add-Some-iff*)

apply *auto*

done

lemma *chg-map-new* [*simp*]:

$m \ a = \text{None} \implies \text{chg-map } f \ a \ m = m$

by (*unfold chg-map-def, auto*)

lemma *chg-map-upd* [*simp*]:

$m \ a = \text{Some } b \implies \text{chg-map } f \ a \ m = m(a | \rightarrow f \ b)$

by (*unfold chg-map-def, auto*)

lemma *ran-override-chg-map*:

$A \in \text{dom } G \implies \text{ran } (G ++ \text{Map.empty}(A | \rightarrow B)) = (\text{ran } (\text{chg-map } (\lambda x. B) \ A \ G))$

apply (*unfold dom-def ran-def*)

apply (*subst map-add-Some-iff [THEN ext]*)

apply *auto*

apply (*rename-tac T*)

apply (*case-tac T = A*)

apply *auto*

done

1.1.6 Part

definition *Part* :: [*'a set, 'b => 'a*] => *'a set* **where**

$\text{Part } A \ h = A \cap \{x. \exists z. x = h(z)\}$

lemma *Part-UNIV-Inl-comp*:

$((\text{Part } \text{UNIV } (\text{Inl } o \ f)) = (\text{Part } \text{UNIV } (\text{Inl } o \ g))) = ((\text{Part } \text{UNIV } \ f) = (\text{Part } \text{UNIV } \ g))$

apply (*unfold Part-def*)

apply *auto*

apply (*erule equalityE*)

apply (*erule subsetCE*)

apply *auto*

apply (*erule equalityE*)

apply (*erule subsetCE*)

apply *auto*

done

lemma *Part-eqI* [*intro*]: $\llbracket a \in A; a=h(b) \rrbracket \implies a \in \text{Part } A \ h$

by (*auto simp add: Part-def*)

lemmas *PartI = Part-eqI [OF - refl]*

lemma *PartE [elim!]*: $\llbracket a \in \text{Part } A \ h; \ !z. \llbracket a \in A; \ a=h(z) \rrbracket \implies P \rrbracket \implies P$
by (*auto simp add: Part-def*)

lemma *Part-subset*: $\text{Part } A \ h \subseteq A$
by (*auto simp add: Part-def*)

lemma *Part-mono*: $A \subseteq B \implies \text{Part } A \ h \subseteq \text{Part } B \ h$
by *blast*

lemmas *basic-monos = basic-monos Part-mono*

lemma *PartD1*: $a \in \text{Part } A \ h \implies a \in A$
by (*simp add: Part-def*)

lemma *Part-id*: $\text{Part } A \ (\lambda x. x) = A$
by *blast*

lemma *Part-Int*: $\text{Part } (A \cap B) \ h = (\text{Part } A \ h) \cap (\text{Part } B \ h)$
by *blast*

lemma *Part-Collect*: $\text{Part } (A \cap \{x. P \ x\}) \ h = (\text{Part } A \ h) \cap \{x. P \ x\}$
by *blast*

1.1.7 Set operators

lemma *subset-lemma*:
 $\llbracket A \cap B = \{\}; \ A \subseteq B \rrbracket \implies A = \{\}$
by *auto*

lemma *subset-lemma2*:
 $\llbracket B \cap A = \{\}; \ C \subseteq A \rrbracket \implies C \cap B = \{\}$
by *auto*

lemma *insert-inter*:
 $\llbracket a \notin A; \ (A \cap B) = \{\} \rrbracket \implies (A \cap (\text{insert } a \ B)) = \{\}$
by *auto*

lemma *insert-notmem*:
 $\llbracket a \neq b; \ a \notin B \rrbracket \implies a \notin (\text{insert } b \ B)$
by *auto*

lemma *insert-union*:
 $A \cup (\text{insert } a \ B) = \text{insert } a \ A \cup B$
by *auto*

lemma *insert-or*:

$\{s. s = t1 \vee (P s)\} = \text{insert } t1 \{s. P s\}$
by auto

lemma *Collect-subset*:

$\{x. x \subseteq A \wedge P x\} = \{x \in \text{Pow } A. P x\}$
by auto

lemma *OneElement-Card* [simp]:

$\llbracket \text{finite } M; \text{card } M \leq \text{Suc } 0; t \in M \rrbracket \implies M = \{t\}$
apply auto
apply (*rename-tac s*)
apply (*rule-tac P=finite M in mp*)
apply (*rule-tac P=card M ≤ Suc 0 in mp*)
apply (*rule-tac P=t ∈ M in mp*)
apply (*rule-tac F=M in finite-induct*)
apply auto
apply (*rule-tac P=finite M in mp*)
apply (*rule-tac P=card M ≤ Suc 0 in mp*)
apply (*rule-tac P=s ∈ M in mp*)
apply (*rule-tac P=t ∈ M in mp*)
apply (*rule-tac F=M in finite-induct*)
apply auto
done

1.1.8 One point rule

lemma *Ex1-one-point* [simp]:

$(\exists! x. P x \wedge x = a) = P a$
by auto

lemma *Ex1-one-point2* [simp]:

$(\exists! x. P x \wedge Q x \wedge x = a) = (P a \wedge Q a)$
by auto

lemma *Some-one-point* [simp]:

$P a \implies (\text{SOME } x. P x \wedge x = a) = a$
by auto

lemma *Some-one-point2* [simp]:

$\llbracket Q a; P a \rrbracket \implies (\text{SOME } x. P x \wedge Q x \wedge x = a) = a$
by auto

end

2 Partitoned Data Spaces for Statecharts

theory *DataSpace*
imports *Contrib*
begin

2.1 Definitions

definition

```
DataSpace :: ('d set) list
            => bool where
DataSpace L = ((distinct L) ∧
               (∀ D1∈(set L). ∀ D2∈(set L).
                D1 ≠ D2 → ((D1 ∩ D2) = {})) ∧
               ((∪ (set L)) = UNIV))
```

lemma *DataSpace-EmptySet*:

```
[UNIV] ∈ { L | L. DataSpace L }
by (unfold DataSpace-def, auto)
```

definition *dataspace* = { L | (L::('d set) list). DataSpace L }

typedef 'd *dataspace* = *dataspace* :: 'd set list set

```
unfolding dataspace-def
apply (rule exI)
apply (rule DataSpace-EmptySet)
done
```

definition

```
PartNum :: ('d)dataspace => nat where
PartNum = length o Rep-dataspace
```

definition

```
PartDom :: ['d dataspace, nat] => ('d set) (infixl !D! 101) where
PartDom d n = (Rep-dataspace d) ! n
```

2.2 Lemmas

2.2.1 *DataSpace*

lemma *DataSpace-UNIV* [*simp*]:

```
DataSpace [UNIV]
by (unfold DataSpace-def, auto)
```

lemma *DataSpace-select*:

```
DataSpace (Rep-dataspace L)
apply (cut-tac x=L in Rep-dataspace)
apply (unfold dataspace-def)
apply auto
done
```

lemma *UNIV-dataspace* [*simp*]:

```
[UNIV] ∈ dataspace
by (unfold dataspace-def, auto)
```

lemma *Inl-Inr-DataSpace* [*simp*]:

```

  DataSpace [Part UNIV Inl, Part UNIV Inr]
apply (unfold DataSpace-def)
apply auto
apply (rename-tac d)
apply (rule-tac b=(inv Inl) d in Part-eqI)
apply auto
apply (rule sym)
apply (case-tac d)
apply auto
done

```

```

lemma Inl-Inr-dataspace [simp]:
  [Part UNIV Inl, Part UNIV Inr] ∈ dataspace
by (unfold dataspace-def, auto)

```

```

lemma InlInr-InlInl-Inr-DataSpace [simp]:
  DataSpace [Part UNIV (Inl o Inr), Part UNIV (Inl o Inl), Part UNIV Inr]
apply (unfold DataSpace-def)
apply auto
apply (unfold Part-def)
apply auto
apply (rename-tac x)
apply (case-tac x)
apply auto
apply (rename-tac a)
apply (case-tac a)
apply auto
done

```

```

lemma InlInr-InlInl-Inr-dataspace [simp]:
  [Part UNIV (Inl o Inr), Part UNIV (Inl o Inl), Part UNIV Inr] : dataspace
by (unfold dataspace-def, auto)

```

2.2.2 *PartNum*

```

lemma PartDom-PartNum-distinct:
   $\llbracket i < \text{PartNum } d; j < \text{PartNum } d; \\ i \neq j; p \in (d \text{ !}D! i) \rrbracket \implies \\ p \notin (d \text{ !}D! j)$ 
apply auto
apply (cut-tac L=d in DataSpace-select)
apply (unfold DataSpace-def)
apply auto
apply (erule-tac x=Rep-dataspace d ! i in ballE)
apply (erule-tac x=Rep-dataspace d ! j in ballE)
apply auto
apply (simp add:distinct-conv-nth PartNum-def)
apply (unfold PartDom-def PartNum-def)
apply auto

```

done

lemma *PartDom-PartNum-distinct2*:

$\llbracket i < \text{PartNum } d; j < \text{PartNum } d;$
 $i \neq j; p \in (d \text{ !}D! j) \rrbracket \implies$
 $p \notin (d \text{ !}D! i)$

apply *auto*

apply (*cut-tac* $L=d$ **in** *DataSpace-select*)

apply (*unfold* *DataSpace-def*)

apply *auto*

apply (*erule-tac* $x=\text{Rep-dataspace } d \text{ ! } i$ **in** *ballE*)

apply (*erule-tac* $x=\text{Rep-dataspace } d \text{ ! } j$ **in** *ballE*)

apply *auto*

apply (*simp* *add:distinct-conv-nth* *PartNum-def*)

apply (*unfold* *PartDom-def* *PartNum-def*)

apply *auto*

done

lemma *PartNum-length* [*simp*]:

$(\text{DataSpace } L) \implies (\text{PartNum } (\text{Abs-dataspace } L) = (\text{length } L))$

apply (*unfold* *PartNum-def*)

apply *auto*

apply (*subst* *Abs-dataspace-inverse*)

apply (*unfold* *dataspace-def*)

apply *auto*

done

end

3 Data Space Assignments

theory *Data*

imports *DataSpace*

begin

3.1 Total data space assignments

definition

$\text{Data} :: ['d \text{ list}, 'd \text{ dataspace}]$

$\implies \text{bool}$ **where**

$\text{Data } L D = (((\text{length } L) = (\text{PartNum } D)) \wedge$

$(\forall i \in \{n. n < (\text{PartNum } D)\}. (L!i) \in (\text{PartDom } D \ i)))$

lemma *Data-EmptySet*:

$([\text{@ } t. \text{True}], \text{Abs-dataspace } [\text{UNIV}]) \in \{ (L,D) \mid L D. \text{Data } L D \}$

apply (*unfold* *Data-def* *PartDom-def*)

apply *auto*

apply (*subst* *Abs-dataspace-inverse*)

apply *auto*

done

definition

```
data =  
  { (L,D) |  
    (L::('d list))  
    (D::('d dataspace)).  
    Data L D }
```

```
typedef 'd data = data :: ('d list * 'd dataspace) set  
  unfolding data-def  
  apply (rule exI)  
  apply (rule Data-EmptySet)  
  done
```

definition

```
DataValue :: ('d data) => ('d list) where  
DataValue = fst o Rep-data
```

definition

```
DataSpace :: ('d data) => ('d dataspace) where  
DataSpace = snd o Rep-data
```

definition

```
DataPart :: ['d data, nat] => 'd ((- !P!/ -) [10,11]10) where  
DataPart d n = (DataValue d) ! n
```

lemma Rep-data-tuple:

```
Rep-data D = (DataValue D, DataSpace D)  
by (unfold DataValue-def DataSpace-def, simp)
```

lemma Rep-data-select:

```
(DataValue D, DataSpace D) ∈ data  
apply (subst Rep-data-tuple [THEN sym])  
apply (rule Rep-data)  
done
```

lemma Data-select:

```
Data (DataValue D) (DataSpace D)  
apply (cut-tac D=D in Rep-data-select)  
apply (unfold data-def)  
apply auto  
done
```

lemma length-DataValue-PartNum [simp]:

```
length (DataValue D) = PartNum (Data.DataSpace D)  
apply (cut-tac D=D in Data-select)  
apply (unfold Data-def)  
apply auto
```

done

lemma *DataValue-PartDom* [simp]:

$i < \text{PartNum } (\text{Data.DataSpace } D) \implies$

$\text{DataValue } D ! i \in \text{PartDom } (\text{Data.DataSpace } D) \ i$

apply (*cut-tac D=D in Data-select*)

apply (*unfold Data-def*)

apply *auto*

done

lemma *DataPart-PartDom* [simp]:

$i < \text{PartNum } (\text{Data.DataSpace } d) \longrightarrow (d !P! i) \in ((\text{Data.DataSpace } d) !D! i)$

apply (*unfold DataPart-def*)

apply *auto*

done

3.2 Partial data space assignments

definition

$PData :: ['d \text{ option list}, 'd \text{ dataspace}] \Rightarrow \text{bool}$ **where**

$PData \ L \ D == ((\text{length } L) = (\text{PartNum } D)) \wedge$

$(\forall i \in \{n. n < (\text{PartNum } D)\}.$

$(L!i) \neq \text{None} \longrightarrow \text{the } (L!i) \in (\text{PartDom } D \ i))$

lemma *PData-EmptySet*:

$([\text{Some } (@ \ t. \ \text{True})], \text{Abs-dataspace } [UNIV]) \in \{ (L,D) \mid L \ D. \ PData \ L \ D \}$

apply (*unfold PData-def PartDom-def*)

apply *auto*

apply (*subst Abs-dataspace-inverse*)

apply *auto*

done

definition

$pdata =$

$\{ (L,D) \mid$

$(L::('d \text{ option list}))$

$(D::('d \text{ dataspace})).$

$PData \ L \ D \}$

typedef $'d \ pdata = pdata :: ('d \ \text{option list} * 'd \ \text{dataspace}) \ \text{set}$

unfolding *pdata-def*

apply (*rule exI*)

apply (*rule PData-EmptySet*)

done

definition

$PDataValue :: ('d \ pdata) \Rightarrow ('d \ \text{option list})$ **where**

$PDataValue = \text{fst } o \ \text{Rep-pdata}$

definition

$PDataSpace :: ('d\ pdata) => ('d\ dataspace)$ **where**
 $PDataSpace = snd\ o\ Rep-pdata$

definition

$Data2PData :: ('d\ data) => ('d\ pdata)$ **where**
 $Data2PData\ D = (let$
 $(L,DP) = Rep-data\ D;$
 $OL = map\ Some\ L$
in
 $Abs-pdata\ (OL,DP))$

definition

$PData2Data :: ('d\ pdata) => ('d\ data)$ **where**
 $PData2Data\ D = (let$
 $(OL,DP) = Rep-pdata\ D;$
 $L = map\ the\ OL$
in
 $Abs-data\ (L,DP))$

definition

$DefaultPData :: ('d\ dataspace) => ('d\ pdata)$ **where**
 $DefaultPData\ D = Abs-pdata\ (replicate\ (PartNum\ D)\ None,\ D)$

definition

$OptionOverride :: ('d\ option * 'd) => 'd$ **where**
 $OptionOverride\ P = (if\ (fst\ P) = None\ then\ (snd\ P)\ else\ (the\ (fst\ P)))$

definition

$DataOverride :: ['d\ pdata,\ 'd\ data] => 'd\ data\ ((-\ [D+]/ -)\ [10,11]10)$ **where**
 $DataOverride\ D1\ D2 =$
 $(let$
 $(L1,DP1) = Rep-pdata\ D1;$
 $(L2,DP2) = Rep-pdata\ D2;$
 $L = map\ OptionOverride\ (zip\ L1\ L2)$
 in
 $Abs-data\ (L,DP2))$

lemma Rep-pdata-tuple:

$Rep-pdata\ D = (PDataValue\ D,\ PDataSpace\ D)$
apply $(unfold\ PData\ Value-def\ PDataSpace-def)$
apply $(simp)$
done

lemma Rep-pdata-select:

$(PDataValue\ D,\ PDataSpace\ D) \in pdata$
apply $(subst\ Rep-pdata-tuple\ [THEN\ sym])$
apply $(rule\ Rep-pdata)$
done

lemma *PData-select*:
 $PData (PDataValue D) (PDataSpace D)$
apply (*cut-tac* $D=D$ **in** *Rep-pdata-select*)
apply (*unfold* *pdata-def*)
apply *auto*
done

3.2.1 *DefaultPData*

lemma *PData-DefaultPData* [*simp*]:
 $PData (replicate (PartNum D) None) D$
apply (*unfold* *PData-def*)
apply *auto*
done

lemma *pdata-DefaultPData* [*simp*]:
 $(replicate (PartNum D) None, D) \in pdata$
apply (*unfold* *pdata-def*)
apply *auto*
done

lemma *PDataSpace-DefaultPData* [*simp*]:
 $PDataSpace (DefaultPData D) = D$
apply (*unfold* *DataSpace-def* *PDataSpace-def* *DefaultPData-def*)
apply *auto*
apply (*subst* *Abs-pdata-inverse*)
apply *auto*
done

lemma *length-PartNum-PData* [*simp*]:
 $length (PDataValue P) = PartNum (PDataSpace P)$
apply (*cut-tac* $D=P$ **in** *Rep-pdata-select*)
apply (*unfold* *pdata-def* *PData-def*)
apply *auto*
done

3.2.2 *Data2PData*

lemma *PData-Data2PData* [*simp*]:
 $PData (map Some (DataValue D)) (Data.DataSpace D)$
apply (*unfold* *PData-def*)
apply *auto*
done

lemma *pdata-Data2PData* [*simp*]:
 $(map Some (DataValue D), Data.DataSpace D) \in pdata$
apply (*unfold* *pdata-def*)
apply *auto*
done

```

lemma DataSpace-Data2PData [simp]:
  (PDataSpace (Data2PData D)) = (Data.DataSpace D)
apply (unfold DataSpace-def PDataSpace-def Data2PData-def Let-def)
apply auto
apply (cut-tac D=D in Rep-data-tuple)
apply auto
apply (subst Abs-pdata-inverse)
apply auto
done

```

```

lemma PDataValue-Data2PData-DataValue [simp]:
  (map the (PDataValue (Data2PData D))) = DataValue D
apply (unfold DataValue-def PDataValue-def Data2PData-def Let-def)
apply auto
apply (cut-tac D=D in Rep-data-tuple)
apply auto
apply (subst Abs-pdata-inverse)
apply simp
apply (simp del: map-map)
done

```

```

lemma DataSpace-PData2Data:
  Data (map the (PDataValue D)) (PDataSpace D) ==>
  (Data.DataSpace (PData2Data D) = (PDataSpace D))
apply (unfold DataSpace-def PDataSpace-def PData2Data-def Let-def)
apply auto
apply (cut-tac D=D in Rep-pdata-tuple)
apply auto
apply (subst Abs-data-inverse)
apply (unfold data-def)
apply auto
done

```

```

lemma PartNum-PDataValue-PartDom [simp]:
  [[ i < PartNum (PDataSpace Q);
    PDataValue Q ! i = Some y ]] ==>
  y ∈ PartDom (PDataSpace Q) i
apply (cut-tac D=Q in Rep-pdata-select)
apply (unfold pdata-def PData-def)
apply auto
done

```

3.2.3 DataOverride

```

lemma Data-DataOverride:
  ((PDataSpace P) = (Data.DataSpace Q)) ==>
  Data (map OptionOverride (zip (PDataValue P) (Data.DataValue Q))) (Data.DataSpace Q)

```

```

apply (unfold Data-def)
apply auto
apply (unfold OptionOverride-def)
apply auto
apply (rename-tac i D)
apply (case-tac PDataValue P ! i = None)
apply auto
apply (drule sym)
apply auto
done

```

```

lemma data-DataOverride:
  ((PDataSpace P) = (Data.DataSpace Q))  $\implies$ 
  (map OptionOverride (zip (PDataValue P) (Data.DataValue Q)), Data.DataSpace
  Q)  $\in$  data
apply (unfold data-def)
apply auto
apply (rule Data-DataOverride)
apply fast
done

```

```

lemma DataSpace-DataOverride [simp]:
  ((Data.DataSpace D) = (PDataSpace E))  $\implies$ 
  Data.DataSpace (E [D+] D) = (Data.DataSpace D)
apply (unfold DataSpace-def DataOverride-def Let-def)
apply auto
apply (cut-tac D=D in Rep-data-tuple)
apply (cut-tac D=E in Rep-pdata-tuple)
apply auto
apply (subst Abs-data-inverse)
apply auto
apply (drule sym)
apply simp
apply (rule data-DataOverride)
apply auto
done

```

```

lemma DataValue-DataOverride [simp]:
  ((PDataSpace P) = (Data.DataSpace Q))  $\implies$ 
  (DataValue (P [D+] Q) = (map OptionOverride (zip (PDataValue P) (Data.DataValue
  Q))))
apply (unfold DataValue-def DataOverride-def Let-def)
apply auto
apply (cut-tac D=P in Rep-pdata-tuple)
apply (cut-tac D=Q in Rep-data-tuple)
apply auto
apply (subst Abs-data-inverse)
apply auto
apply (rule data-DataOverride)

```

```

apply auto
done

```

3.2.4 OptionOverride

```

lemma DataValue-OptionOverride-nth:
   $\llbracket ((PDataSpace\ P) = (DataSpace\ Q));$ 
   $i < PartNum\ (DataSpace\ Q) \rrbracket \implies$ 
   $(DataValue\ (P\ [D+] \ Q) ! i) =$ 
   $OptionOverride\ (PDataValue\ P\ !\ i,\ DataValue\ Q\ !\ i)$ 
apply auto
done

```

```

lemma None-OptionOverride [simp]:
   $(fst\ P) = None \implies OptionOverride\ P = (snd\ P)$ 
apply (unfold OptionOverride-def)
apply auto
done

```

```

lemma Some-OptionOverride [simp]:
   $(fst\ P) \neq None \implies OptionOverride\ P = the\ (fst\ P)$ 
apply (unfold OptionOverride-def)
apply auto
done

```

```

end

```

4 Update-Functions on Data Spaces

```

theory Update
imports Data
begin

```

4.1 Total update-functions

```

definition
  Update :: (('d data) => ('d data)) => bool where
  Update U = ( $\forall d. Data.DataSpace\ d = DataSpace\ (U\ d)$ )

```

```

lemma Update-EmptySet:
   $(\% d. d) \in \{ L \mid L. Update\ L \}$ 
by (unfold Update-def, auto)

```

```

definition
  update =  $\{ L \mid (L::('d\ data) \Rightarrow ('d\ data))). Update\ L \}$ 

```

```

typedef 'd update = update :: ('d data => 'd data) set
unfolding update-def
apply (rule exI)

```

apply (*rule Update-EmptySet*)
done

definition

UpdateApply :: [*d update*, *d data*] => *d data* ((- !!!/ -) [10,11]10) **where**
UpdateApply *U D* == *Rep-update U D*

definition

DefaultUpdate :: (*d update*) **where**
DefaultUpdate == *Abs-update* ($\lambda D. D$)

4.1.1 Basic lemmas

lemma *Update-select*:

Update (*Rep-update U*)
apply (*cut-tac x=U in Rep-update*)
apply (*unfold update-def*)
apply *auto*
done

lemma *DataSpace-DataSpace-Update* [*simp*]:

Data.DataSpace (*Rep-update U DP*) = *Data.DataSpace DP*
apply (*cut-tac U=U in Update-select*)
apply (*unfold Update-def*)
apply *auto*
done

4.1.2 DefaultUpdate

lemma *Update-DefaultUpdate* [*simp*]:

Update ($\lambda D. D$)
by (*unfold Update-def, auto*)

lemma *update-DefaultUpdate* [*simp*]:

$(\lambda D. D) \in \text{update}$
by (*unfold update-def, auto*)

lemma *DataSpace-UpdateApply* [*simp*]:

Data.DataSpace (*U !!! D*) = *Data.DataSpace D*
by (*unfold UpdateApply-def, auto*)

4.2 Partial update-functions

definition

PUpdate :: ((*d data*) => (*d pdata*)) => *bool* **where**
PUpdate U = ($\forall d. \text{Data.DataSpace } d = \text{PDataSpace } (U d)$)

lemma *PUpdate-EmptySet*:

$(\% d. \text{Data2PData } d) \in \{ L \mid L. \text{PUpdate } L \}$
by (*unfold PUpdate-def, auto*)

definition $pupdate = \{ L \mid (L::('d\ data) \Rightarrow ('d\ pdata))). PUpdate\ L\}$

typedef $'d\ pupdate = pupdate :: ('d\ data \Rightarrow 'd\ pdata)\ set$
unfolding $pupdate-def$
apply $(rule\ exI)$
apply $(rule\ PUpdate-EmptySet)$
done

definition

$PUpdateApply :: ['d\ pupdate, 'd\ data] \Rightarrow 'd\ pdata\ ((- \ \!/\ -)\ [10,11]10)$ **where**
 $PUpdateApply\ U\ D = Rep-pupdate\ U\ D$

definition

$DefaultPUpdate :: ('d\ pupdate)$ **where**
 $DefaultPUpdate = Abs-pupdate\ (\lambda\ D.\ DefaultPData\ (Data.DataSpace\ D))$

4.2.1 Basic lemmas

lemma $PUpdate-select$:

$PUpdate\ (Rep-pupdate\ U)$
apply $(cut-tac\ x=U\ \mathbf{in}\ Rep-pupdate)$
apply $(unfold\ pupdate-def)$
apply $auto$
done

lemma $DataSpace-PDataSpace-PUpdate\ [simp]$:

$PDataSpace\ (Rep-pupdate\ U\ DP) = Data.DataSpace\ DP$
apply $(cut-tac\ U=U\ \mathbf{in}\ PUpdate-select)$
apply $(unfold\ PUpdate-def)$
apply $auto$
done

4.2.2 Data2PData

lemma $PUpdate-Data2PData\ [simp]$:

$PUpdate\ Data2PData$
by $(unfold\ PUpdate-def,\ auto)$

lemma $pupdate-Data2PData\ [simp]$:

$Data2PData \in pupdate$
by $(unfold\ pupdate-def,\ auto)$

4.2.3 PUpdate

lemma $PUpdate-DefaultPUpdate\ [simp]$:

$PUpdate\ (\lambda\ D.\ DefaultPData\ (Data.DataSpace\ D))$
apply $(unfold\ PUpdate-def)$
apply $auto$
done

lemma *pupdate-DefaultPUpdate* [*simp*]:
 $(\lambda D. \text{DefaultPData } (Data.DataSpace D)) \in \text{pupdate}$
apply (*unfold pupdate-def*)
apply *auto*
done

lemma *DefaultPUpdate-None* [*simp*]:
 $(\text{DefaultPUpdate} !! D) = \text{DefaultPData } (DataSpace D)$
apply (*unfold DefaultPUpdate-def PUpdateApply-def*)
apply (*subst Abs-pupdate-inverse*)
apply *auto*
done

4.2.4 SequentialRacing

definition

UpdateOverride :: [*d pupdate*, *d update*] =>
 $'d \text{ update } ((- [U+]/ -) [10,11]10) \text{ where}$
 $\text{UpdateOverride } U P = \text{Abs-update } (\lambda DA . (U !! DA) [D+] (P !!! DA))$

inductive

FoldSet :: (*'b* => *'a* => *'a*) => *'a* => *'b set* => *'a* => *bool*
for *h* :: *'b* => *'a* => *'a*
and *z* :: *'a*

where

emptyI [*intro*]: *FoldSet h z* {} *z*
| *insertI* [*intro*]:
 $\llbracket x \notin A; \text{FoldSet } h z A y \rrbracket$
 $\implies \text{FoldSet } h z (\text{insert } x A) (h x y)$

definition

SequentialRacing :: (*d pupdate set*) => (*d update set*) **where**
 $\text{SequentialRacing } U =$
 $\{u. \text{FoldSet } \text{UpdateOverride } \text{DefaultUpdate } U u\}$

lemma

FoldSet-imp-finite:

$\text{FoldSet } h z A x \implies \text{finite } A$

by (*induct set: FoldSet*) *auto*

lemma

finite-imp-FoldSet:

$\text{finite } A \implies \exists x. \text{FoldSet } h z A x$

by (induct set: finite) auto

lemma *finite-SequentialRacing*:

$finite\ US \implies (SOME\ u.\ u \in SequentialRacing\ US) \in SequentialRacing\ US$

apply (unfold SequentialRacing-def)

apply auto

apply (drule-tac h=UpdateOverride and z=DefaultUpdate in finite-imp-FoldSet)

apply auto

apply (rule someI)

apply auto

done

end

5 Label Expressions

theory *Expr*

imports *Update*

begin

datatype $(s,e)expr = true$

| *In* s

| *En* e

| *NOT* $(s,e)expr$

| *And* $(s,e)expr\ (s,e)expr$

| *Or* $(s,e)expr\ (s,e)expr$

type-synonym $'d\ guard = ('d\ data) \implies bool$

type-synonym $(e,d)action = (e\ set * 'd\ pupdate)$

type-synonym $(s,e,'d)label = ((s,e)expr * 'd\ guard * (e,d)action)$

type-synonym $(s,e,'d)trans = (s * (s,e,'d)label * 's)$

primrec

$eval_expr :: [(s\ set * 'e\ set), (s,e)expr] \implies bool$ **where**

$eval_expr\ sc\ true = True$

| $eval_expr\ sc\ (En\ ev) = (ev \in snd\ sc)$

| $eval_expr\ sc\ (In\ st) = (st \in fst\ sc)$

| $eval_expr\ sc\ (NOT\ e1) = (\neg (eval_expr\ sc\ e1))$

| $eval_expr\ sc\ (And\ e1\ e2) = ((eval_expr\ sc\ e1) \wedge (eval_expr\ sc\ e2))$

| $eval_expr\ sc\ (Or\ e1\ e2) = ((eval_expr\ sc\ e1) \vee (eval_expr\ sc\ e2))$

primrec

$ExprEvents :: (s,e)expr \implies 'e\ set$ **where**

$ExprEvents\ true = \{\}$

| $ExprEvents\ (En\ ev) = \{ev\}$

| $ExprEvents\ (In\ st) = \{\}$

| $ExprEvents\ (NOT\ e) = (ExprEvents\ e)$

| $ExprEvents\ (And\ e1\ e2) = ((ExprEvents\ e1) \cup (ExprEvents\ e2))$

| $ExprEvents (Or\ e1\ e2) = ((ExprEvents\ e1) \cup (ExprEvents\ e2))$

datatype ($'s, 'e, dead\ 'd$)*atomar* =
 TRUE
 | *FALSE*
 | *IN 's*
 | *EN 'e*
 | *VAL 'd data => bool*

definition
source :: ($'s, 'e, 'd$)*trans* => $'s$ **where**
source t = *fst t*

definition
Source :: ($'s, 'e, 'd$)*trans set* => $'s$ *set* **where**
Source T == *source ' T*

definition
target :: ($'s, 'e, 'd$)*trans* => $'s$ **where**
target t = *snd(snd t)*

definition
Target :: ($'s, 'e, 'd$)*trans set* => $'s$ *set* **where**
Target T = *target ' T*

definition
label :: ($'s, 'e, 'd$)*trans* => ($'s, 'e, 'd$)*label* **where**
label t = *fst (snd t)*

definition
Label :: ($'s, 'e, 'd$)*trans set* => ($'s, 'e, 'd$)*label set* **where**
Label T = *label ' T*

definition
expr :: ($'s, 'e, 'd$)*label* => ($'s, 'e$)*expr* **where**
expr = *fst*

definition
action :: ($'s, 'e, 'd$)*label* => ($'e, 'd$)*action* **where**
action = *snd o snd*

definition
Action :: ($'s, 'e, 'd$)*label set* => ($'e, 'd$)*action set* **where**
Action L = *action ' L*

definition
pupdate :: ($'s, 'e, 'd$)*label* => $'d$ *pupdate* **where**

lemma *Target-EmptySet* [simp]:
 $Target \{\} = \{\}$
by (*unfold Target-def, auto*)

lemma *Label-EmptySet* [simp]:
 $Label \{\} = \{\}$
by (*unfold Label-def, auto*)

lemma *Action-EmptySet* [simp]:
 $Action \{\} = \{\}$
by (*unfold Action-def, auto*)

lemma *PUpdate-EmptySet* [simp]:
 $PUpdate \{\} = \{\}$
by (*unfold PUpdate-def, auto*)

lemma *Actevent-EmptySet* [simp]:
 $Actevent \{\} = \{\}$
by (*unfold Actevent-def, auto*)

lemma *Union-Actevent-subset*:
 $\llbracket m \in M; ((\bigcup (Actevent (Label (Union M)))) \subseteq (N::'a set)) \rrbracket \implies$
 $((\bigcup (Actevent (Label m))) \subseteq N)$
by (*unfold Actevent-def Label-def, auto*)

lemma *action-select* [simp]:
 $action (a,b,c) = c$
by (*unfold action-def, auto*)

lemma *label-select* [simp]:
 $label (a,b,c) = b$
by (*unfold label-def, auto*)

lemma *target-select* [simp]:
 $target (a,b,c) = c$
by (*unfold target-def, auto*)

lemma *actevent-select* [simp]:
 $actevent (a,b,(c,d)) = c$
by (*unfold actevent-def, auto*)

lemma *pupdate-select* [simp]:
 $pupdate (a,b,c,d) = d$
by (*unfold pupdate-def, auto*)

lemma *source-select* [simp]:
 $source (a,b) = a$
by (*unfold source-def, auto*)

```

lemma finite-PUupdate [simp]:
  finite S  $\implies$  finite(PUpdate S)
by (unfold PUpdate-def, auto)

lemma finite-Label [simp]:
  finite S  $\implies$  finite(Label S)
by (unfold Label-def, auto)

lemma fst-defaultaction [simp]:
  fst defaultaction = {}
by (unfold defaultaction-def, auto)

lemma action-defaultlabel [simp]:
  (action defaultlabel) = defaultaction
by (unfold defaultlabel-def action-def, auto)

lemma fst-defaultlabel [simp]:
  (fst defaultlabel) = defaultexpr
by (unfold defaultlabel-def, auto)

lemma ExprEvents-defaultexpr [simp]:
  (ExprEvents defaultexpr) = {}
by (unfold defaultexpr-def, auto)

lemma defaultlabel-defaultexpr [simp]:
  expr defaultlabel = defaultexpr
by (unfold defaultlabel-def expr-def, auto)

lemma target-Target [simp]:
   $t \in T \implies$  target t  $\in$  Target T
by (unfold Target-def, auto)

lemma Source-union : Source s  $\cup$  Source t = Source (s  $\cup$  t)
apply (unfold Source-def)
apply auto
done

lemma Target-union : Target s  $\cup$  Target t = Target (s  $\cup$  t)
apply (unfold Target-def)
apply auto
done

end

```

6 Sequential Automata

```

theory SA
imports Expr
begin

```

definition

$SeqAuto :: ['s\ set,$
 $\quad 's,$
 $\quad (('s,'e,'d)label)\ set,$
 $\quad (('s,'e,'d)trans)\ set]$
 $=> bool\ \mathbf{where}$
 $SeqAuto\ S\ I\ L\ D = (I \in S \wedge S \neq \{\}) \wedge finite\ S \wedge finite\ D \wedge$
 $(\forall (s,l,t) \in D. s \in S \wedge t \in S \wedge l \in L)$

lemma *SeqAuto-EmptySet:*

$(\{ @x . True \}, (\@x . True), \{\}, \{\}) \in \{(S,I,L,D) \mid S\ I\ L\ D. SeqAuto\ S\ I\ L\ D\}$
by (*unfold SeqAuto-def, auto*)

definition

$seqauto =$
 $\{ (S,I,L,D) \mid$
 $\quad (S::'s\ set)$
 $\quad (I::'s)$
 $\quad (L::('s,'e,'d)label)\ set)$
 $\quad (D::('s,'e,'d)trans)\ set).$
 $SeqAuto\ S\ I\ L\ D\}$

typedef $('s,'e,'d)\ seqauto =$

$seqauto :: ('s\ set * 's * (('s,'e,'d)label)\ set * (('s,'e,'d)trans)\ set)\ set$
unfolding *seqauto-def*
apply (*rule exI*)
apply (*rule SeqAuto-EmptySet*)
done

definition

$States :: (('s,'e,'d)seqauto) => 's\ set\ \mathbf{where}$
 $States = fst\ o\ Rep-seqauto$

definition

$InitState :: (('s,'e,'d)seqauto) => 's\ \mathbf{where}$
 $InitState = fst\ o\ snd\ o\ Rep-seqauto$

definition

$Labels :: (('s,'e,'d)seqauto) => (('s,'e,'d)label)\ set\ \mathbf{where}$
 $Labels = fst\ o\ snd\ o\ snd\ o\ Rep-seqauto$

definition

$Delta :: (('s,'e,'d)seqauto) => (('s,'e,'d)trans)\ set\ \mathbf{where}$
 $Delta = snd\ o\ snd\ o\ snd\ o\ Rep-seqauto$

definition

$SAEvents :: (('s,'e,'d)seqauto) => 'e\ set\ \mathbf{where}$
 $SAEvents\ SA = (\bigcup l \in Label\ (Delta\ SA). (fst\ (action\ l)) \cup (ExprEvents\ (expr\ l)))$

lemma *Rep-seqauto-tuple*:
 $Rep-seqauto\ SA = (States\ SA, InitState\ SA, Labels\ SA, Delta\ SA)$
by (*unfold States-def InitState-def Labels-def Delta-def, auto*)

lemma *Rep-seqauto-select*:
 $(States\ SA, InitState\ SA, Labels\ SA, Delta\ SA) \in seqauto$
by (*rule Rep-seqauto-tuple [THEN subst], rule Rep-seqauto*)

lemma *SeqAuto-select*:
 $SeqAuto\ (States\ SA)\ (InitState\ SA)\ (Labels\ SA)\ (Delta\ SA)$
by (*cut-tac SA=SA in Rep-seqauto-select, unfold seqauto-def, auto*)

lemma *neq-States [simp]*:
 $States\ SA \neq \{\}$
apply (*cut-tac Rep-seqauto-select*)
apply *auto*
apply (*unfold seqauto-def SeqAuto-def*)
apply *auto*
done

lemma *SA-States-disjunct* :
 $(States\ A) \cap (States\ A') = \{\} \implies A' \neq A$
by *auto*

lemma *SA-States-disjunct2* :
 $\llbracket (States\ A) \cap C = \{\}; States\ B \subseteq C \rrbracket \implies B \neq A$
apply (*rule SA-States-disjunct*)
apply *auto*
done

lemma *SA-States-disjunct3* :
 $\llbracket C \cap States\ A = \{\}; States\ B \subseteq C \rrbracket \implies States\ A \cap States\ B = \{\}$
apply (*cut-tac SA=B in neq-States*)
apply *fast*
done

lemma *EX-State-SA [simp]*:
 $\exists S. S \in States\ SA$
apply (*cut-tac Rep-seqauto-select*)
apply (*unfold seqauto-def SeqAuto-def*)
apply *auto*
done

lemma *finite-States [simp]*:
 $finite\ (States\ A)$
apply (*cut-tac Rep-seqauto-select*)
apply (*unfold seqauto-def SeqAuto-def*)
apply *auto*

done

lemma *finite-Delta* [*simp*]:
 finite (*Delta A*)
apply (*cut-tac Rep-seqauto-select*)
apply (*unfold seqauto-def SeqAuto-def*)
apply *auto*
done

lemma *InitState-States* [*simp*]:
 InitState A \in *States A*
apply (*cut-tac Rep-seqauto-select*)
apply (*unfold seqauto-def SeqAuto-def*)
apply *auto*
done

lemma *SeqAuto-EmptySet-States* [*simp*]:
 (*States* (*Abs-seqauto* ($\{\text{@x. True}\}$, (@x. True), $\{\}$, $\{\}$))) = $\{\text{@x. True}\}$
apply (*unfold States-def*)
apply (*simp*)
apply (*subst Abs-seqauto-inverse*)
apply (*unfold seqauto-def*)
apply (*rule SeqAuto-EmptySet*)
apply *auto*
done

lemma *SeqAuto-EmptySet-SAEvents* [*simp*]:
 (*SAEvents* (*Abs-seqauto* ($\{\text{@x. True}\}$, (@x. True), $\{\}$, $\{\}$))) = $\{\}$
apply (*unfold SAEvents-def Delta-def*)
apply *simp*
apply (*subst Abs-seqauto-inverse*)
apply (*unfold seqauto-def*)
apply (*rule SeqAuto-EmptySet*)
apply *auto*
done

lemma *Label-Delta-subset* [*simp*]:
 (*Label* (*Delta SA*)) \subseteq *Labels SA*
apply (*unfold Label-def label-def*)
apply *auto*
apply (*cut-tac SA=SA in SeqAuto-select*)
apply (*unfold SeqAuto-def*)
apply *auto*
done

lemma *Target-SAs-Delta-States*:
 Target (\bigcup (*Delta* ' (*SAs HA*))) \subseteq \bigcup (*States* ' (*SAs HA*))
apply (*unfold image-def Target-def target-def*)
apply *auto*

```

apply (rename-tac SA Source Trigger Guard Action Update Target)
apply (cut-tac SA=SA in SeqAuto-select)
apply (unfold SeqAuto-def)
apply auto
done

```

```

lemma States-Int-not-mem:
   $(\bigcup (\text{States } 'F) \text{ Int States } SA) = \{\} \implies SA \notin F$ 
apply (unfold Int-def)
apply auto
apply (subgoal-tac  $\exists S. S \in \text{States } SA$ )
prefer 2
apply (rule EX-State-SA)
apply (erule exE)
apply (rename-tac T)
apply (erule-tac  $x=T$  in allE)
apply auto
done

```

```

lemma Delta-target-States [simp]:
   $\llbracket T \in \text{Delta } A \rrbracket \implies \text{target } T \in \text{States } A$ 
apply (cut-tac SA=A in SeqAuto-select)
apply (unfold SeqAuto-def source-def target-def)
apply auto
done

```

```

lemma Delta-source-States [simp]:
   $\llbracket T \in \text{Delta } A \rrbracket \implies \text{source } T \in \text{States } A$ 
apply (cut-tac SA=A in SeqAuto-select)
apply (unfold SeqAuto-def source-def target-def)
apply auto
done

```

end

7 Syntax of Hierarchical Automata

```

theory HA
imports SA
begin

```

7.1 Definitions

```

definition
  RootEx ::  $[(('s, 'e, 'd) \text{seqauto}) \text{ set},$ 
     $'s \rightarrow ('s, 'e, 'd) \text{seqauto set}] \Rightarrow \text{bool}$  where
  RootEx F G =  $(\exists! A. A \in F \wedge A \notin \bigcup (\text{ran } G))$ 

```

```

definition

```


$Root :: [((\prime s, \prime e, \prime d)seqauto) set,$
 $\prime s \rightarrow (\prime s, \prime e, \prime d) seqauto set]$
 $=> (\prime s, \prime e, \prime d) seqauto \mathbf{where}$
 $Root F G = (@ A. A \in F \wedge A \notin \bigcup (ran G))$

definition

$MutuallyDistinct :: ((\prime s, \prime e, \prime d)seqauto) set => bool \mathbf{where}$
 $MutuallyDistinct F =$
 $(\forall a \in F. \forall b \in F. a \neq b \rightarrow (States a) \cap (States b) = \{\})$

definition

$OneAncestor :: [((\prime s, \prime e, \prime d)seqauto) set,$
 $\prime s \rightarrow (\prime s, \prime e, \prime d) seqauto set] => bool \mathbf{where}$
 $OneAncestor F G =$
 $(\forall A \in F - \{Root F G\} .$
 $\exists! s. s \in (\bigcup A' \in F - \{A\} . States A') \wedge$
 $A \in the (G s))$

definition

$NoCycles :: [((\prime s, \prime e, \prime d)seqauto) set,$
 $\prime s \rightarrow (\prime s, \prime e, \prime d) seqauto set] => bool \mathbf{where}$
 $NoCycles F G =$
 $(\forall S \in Pow (\bigcup A \in F. States A).$
 $S \neq \{\} \rightarrow (\exists s \in S. S \cap (\bigcup A \in the (G s). States A) = \{\}))$

definition

$IsCompFun :: [((\prime s, \prime e, \prime d)seqauto) set,$
 $\prime s \rightarrow (\prime s, \prime e, \prime d) seqauto set] => bool \mathbf{where}$
 $IsCompFun F G = ((dom G = (\bigcup A \in F. States A)) \wedge$
 $(\bigcup (ran G) = (F - \{Root F G\})) \wedge$
 $(RootEx F G) \wedge$
 $(OneAncestor F G) \wedge$
 $(NoCycles F G))$

7.1.1 Well-formedness for the syntax of HA

definition

$HierAuto :: [\prime d data,$
 $((\prime s, \prime e, \prime d)seqauto) set,$
 $\prime e set,$
 $\prime s \rightarrow ((\prime s, \prime e, \prime d)seqauto) set]$

```

=> bool where
HierAuto D F E G = (( $\bigcup$  A  $\in$  F. SAEvents A)  $\subseteq$  E  $\wedge$ 
  MutuallyDistinct F  $\wedge$ 
  finite F  $\wedge$ 
  IsCompFun F G)

lemma HierAuto-EmptySet:
  ((@x. True), {Abs-seqauto ({@x. True}, (@x. True), {}, {})}, {}), {},
  Map.empty (@x. True  $\mapsto$  {}))  $\in$  {(D,F,E,G) | D F E G. HierAuto D F E G}
apply (unfold HierAuto-def IsCompFun-def Root-def RootEx-def MutuallyDistinct-def
  OneAncestor-def NoCycles-def)
apply auto
done

```

```

definition
hierauto =
  {(D,F,E,G) |
    (D::'d data)
    (F::('s,'e,'d) seqauto) set)
    (E::('e set))
    (G::('s  $\rightarrow$  (('s,'e,'d) seqauto) set)).
    HierAuto D F E G}

```

```

typedef ('s,'e,'d) hierauto =
  hierauto :: ('d data * ('s,'e,'d) seqauto set * 'e set * ('s  $\rightarrow$  ('s,'e,'d) seqauto set))
set
unfolding hierauto-def
apply (rule exI)
apply (rule HierAuto-EmptySet)
done

```

```

definition
SAs :: (('s,'e,'d) hierauto) => (('s,'e,'d) seqauto) set where
SAs = fst o snd o Rep-hierauto

```

```

definition
HAEvents :: (('s,'e,'d) hierauto) => ('e set) where
HAEvents = fst o snd o snd o Rep-hierauto

```

```

definition
CompFun :: (('s,'e,'d) hierauto) => ('s  $\rightarrow$  ('s,'e,'d) seqauto set) where
CompFun = (snd o snd o snd o Rep-hierauto)

```

```

definition
HASates :: (('s,'e,'d) hierauto) => ('s set) where
HASates HA = ( $\bigcup$  A  $\in$  (SAs HA). States A)

```

```

definition
HADelta :: (('s,'e,'d) hierauto) => (('s,'e,'d) trans) set where

```

$HADelta HA = (\bigcup F \in (SAs HA). Delta F)$

definition

$HAINitValue :: (('s, 'e, 'd) hierauto) => 'd \text{ data } \mathbf{where}$
 $HAINitValue == \text{fst } o \text{ Rep-hierauto}$

definition

$HAINitStates :: (('s, 'e, 'd) hierauto) => 's \text{ set } \mathbf{where}$
 $HAINitStates HA == \bigcup A \in (SAs HA). \{ InitState A \}$

definition

$HARoot :: (('s, 'e, 'd) hierauto) => ('s, 'e, 'd) \text{ seqauto } \mathbf{where}$
 $HARoot HA == \text{Root } (SAs HA) (\text{CompFun } HA)$

definition

$HAINitState :: (('s, 'e, 'd) hierauto) => 's \text{ where}$
 $HAINitState HA == \text{InitState } (HARoot HA)$

7.1.2 State successor function

definition

$Chi :: ('s, 'e, 'd) \text{ hierauto} => 's => 's \text{ set } \mathbf{where}$
 $Chi A == (\lambda S \in (HAsStates A) .$
 $\quad \{ S' . \exists SA \in (SAs A) . SA \in \text{the } ((\text{CompFun } A) S) \wedge S' \in \text{States}$
 $SA \})$

definition

$ChiRel :: ('s, 'e, 'd) \text{ hierauto} => ('s * 's) \text{ set } \mathbf{where}$
 $ChiRel A == \{ (S, S') . S \in HAsStates A \wedge S' \in HAsStates A \wedge S' \in (Chi A) S \}$

definition

$ChiPlus :: ('s, 'e, 'd) \text{ hierauto} => ('s * 's) \text{ set } \mathbf{where}$
 $ChiPlus A == (ChiRel A) \hat{+}$

definition

$ChiStar :: ('s, 'e, 'd) \text{ hierauto} => ('s * 's) \text{ set } \mathbf{where}$
 $ChiStar A == (ChiRel A) \hat{*}$

definition

$HigherPriority :: [(('s, 'e, 'd) \text{ hierauto},$
 $\quad ('s, 'e, 'd) \text{ trans} * ('s, 'e, 'd) \text{ trans}] => \text{bool } \mathbf{where}$
 $HigherPriority A ==$
 $\quad \lambda (t, t') \in (HADelta A) \times (HADelta A).$

$$(source\ t', source\ t) \in\ ChiPlus\ A$$

7.1.3 Configurations

definition

$InitConf :: ('s, 'e, 'd)hierauto \Rightarrow 's\ set$ **where**
 $InitConf\ A == (((((HAIInitStates\ A) \times (HAIInitStates\ A)) \cap (ChiRel\ A))^*)$
 $\quad \text{“ } \{HAIInitState\ A\}$

definition

$StepConf :: [('s, 'e, 'd)hierauto, 's\ set,$
 $('s, 'e, 'd)trans\ set] \Rightarrow 's\ set$ **where**

$StepConf\ A\ C\ TS ==$
 $(C - ((ChiStar\ A)\ \text{“}\ (Source\ TS))) \cup$
 $(Target\ TS) \cup$
 $((ChiRel\ A)\ \text{“}\ (Target\ TS)) \cap (HAIInitStates\ A) \cup$
 $((((ChiRel\ A) \cap ((HAIInitStates\ A) \times (HAIInitStates\ A))))^+)$
 $\quad \text{“}\ (((ChiRel\ A)\ \text{“}\ (Target\ TS)) \cap (HAIInitStates\ A))$

7.2 Lemmas

lemma *Rep-hierauto-tuple*:

$Rep-hierauto\ HA = (HAIInitValue\ HA, SAs\ HA, HAEvents\ HA, CompFun\ HA)$

by (*unfold SAs-def HAEvents-def CompFun-def HAIInitValue-def, simp*)

lemma *Rep-hierauto-select*:

$(HAIInitValue\ HA, SAs\ HA, HAEvents\ HA, CompFun\ HA): hierauto$

by (*rule Rep-hierauto-tuple [THEN subst], rule Rep-hierauto*)

lemma *HierAuto-select [simp]*:

$HierAuto\ (HAIInitValue\ HA)\ (SAs\ HA)\ (HAEvents\ HA)\ (CompFun\ HA)$

by (*cut-tac Rep-hierauto-select, unfold hierauto-def, simp*)

7.2.1 *HAStates*

lemma *finite-HAStates* [*simp*]:
 finite (*HAStates HA*)
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold hierauto-def HierAuto-def*)
apply *auto*
apply (*simp add: HAStates-def*)
apply (*rule finite-UN-I*)
apply *fast*
apply (*rule finite-States*)
done

lemma *HAStates-SA-mem*:
 $\llbracket SA \in SAs\ A; S \in States\ SA \rrbracket \implies S \in HAStates\ A$
by (*unfold HAStates-def, auto*)

lemma *ChiRel-HAStates* [*simp*]:
 $(a,b) \in ChiRel\ A \implies a \in HAStates\ A$
apply (*unfold ChiRel-def*)
apply *auto*
done

lemma *ChiRel-HAStates2* [*simp*]:
 $(a,b) \in ChiRel\ A \implies b \in HAStates\ A$
apply (*unfold ChiRel-def*)
apply *auto*
done

7.2.2 *HAEvents*

lemma *HAEvents-SAEvents-SAs*:
 $\bigcup (SAEvents\ ' (SAs\ HA)) \subseteq HAEvents\ HA$
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold hierauto-def HierAuto-def*)
apply *fast*
done

7.2.3 *NoCycles*

lemma *NoCycles-EmptySet* [*simp*]:
 NoCycles $\{\}$ *S*
by (*unfold NoCycles-def, auto*)

lemma *NoCycles-HA* [*simp*]:
 NoCycles (*SAs HA*) (*CompFun HA*)
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold hierauto-def HierAuto-def IsCompFun-def*)
apply *auto*
done

7.2.4 OneAncestor

lemma *OneAncestor-HA* [simp]:
 OneAncestor (SAs HA) (CompFun HA)
apply (cut-tac Rep-hierauto-select)
apply (unfold hierauto-def HierAuto-def IsCompFun-def)
apply auto
done

7.2.5 MutuallyDistinct

lemma *MutuallyDistinct-Single* [simp]:
 MutuallyDistinct {SA}
by (unfold *MutuallyDistinct-def*, auto)

lemma *MutuallyDistinct-EmptySet* [simp]:
 MutuallyDistinct {}
by (unfold *MutuallyDistinct-def*, auto)

lemma *MutuallyDistinct-Insert*:
 [[*MutuallyDistinct* S; (States A) \cap ($\bigcup B \in S.$ States B) = {}]]
 \implies *MutuallyDistinct* (insert A S)
by (unfold *MutuallyDistinct-def*, safe, fast+)

lemma *MutuallyDistinct-Union*:
 [[*MutuallyDistinct* A; *MutuallyDistinct* B;
 ($\bigcup C \in A.$ States C) \cap ($\bigcup C \in B.$ States C) = {}]]
 \implies *MutuallyDistinct* (A \cup B)
by (unfold *MutuallyDistinct-def*, safe, blast+)

lemma *MutuallyDistinct-HA* [simp]:
 MutuallyDistinct (SAs HA)
apply (cut-tac Rep-hierauto-select)
apply (unfold hierauto-def HierAuto-def IsCompFun-def)
apply auto
done

7.2.6 RootEx

lemma *RootEx-Root* [simp]:
 RootEx F G \implies *Root* F G \in F
apply (unfold *RootEx-def* *Root-def*)
apply (erule ex1E)
apply (erule conjE)
apply (rule someI2)
apply blast+
done

lemma *RootEx-Root-ran* [simp]:
 RootEx F G \implies *Root* F G $\notin \bigcup$ (ran G)

```

apply (unfold RootEx-def Root-def)
apply (erule ex1E)
apply (erule conjE)
apply (rule someI2)
apply blast+
done

```

```

lemma RootEx-States-Subset [simp]:
  (RootEx F G)  $\implies$  States (Root F G)  $\subseteq$  ( $\bigcup x \in F . States x$ )
apply (unfold RootEx-def Root-def)
apply (erule ex1E)
apply (erule conjE)
apply (rule someI2)
apply fast
apply (unfold UNION-eq)
apply (simp add: subset-eq)
apply auto
done

```

```

lemma RootEx-States-notdisjunct [simp]:
  RootEx F G  $\implies$  States (Root F G)  $\cap$  ( $\bigcup x \in F . States x$ )  $\neq \{\}$ 
apply (frule RootEx-States-Subset)
apply (case-tac States (Root F G)=\{\})
prefer 2
apply fast
apply simp
done

```

```

lemma Root-neq-SA [simp]:
  [ $\llbracket$  RootEx F G; ( $\bigcup x \in F . States x$ )  $\cap$  States SA = \{\}  $\rrbracket \implies$  Root F G  $\neq$  SA]
apply (rule SA-States-disjunct)
apply (frule RootEx-States-Subset)
apply fast
done

```

```

lemma RootEx-HA [simp]:
  RootEx (SAs HA) (CompFun HA)
apply (cut-tac Rep-hierauto-select)
apply (unfold hierauto-def HierAuto-def IsCompFun-def)
apply fast
done

```

7.2.7 HARoot

```

lemma HARoot-SAs [simp]:
  (HARoot HA)  $\in$  SAs HA
apply (unfold HARoot-def)
apply (cut-tac Rep-hierauto-select)
apply (unfold hierauto-def HierAuto-def)

```

apply *auto*
done

lemma *States-HARoot-HAStates:*
 $States (HARoot HA) \subseteq HAStates HA$
apply (*unfold HAStates-def*)
apply *auto*
apply (*rule-tac x=HARoot HA in bexI*)
apply *auto*
done

lemma *SAEvents-HARoot-HAEvents:*
 $SAEvents (HARoot HA) \subseteq HAEvents HA$
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold hierauto-def HierAuto-def*)
apply *auto*
apply (*rename-tac S*)
apply (*unfold UNION-eq*)
apply (*simp add: subset-eq*)
apply (*erule-tac x=S in alle*)
apply *auto*
done

lemma *HARoot-ran-CompFun:*
 $HARoot HA \notin Union (ran (CompFun HA))$
apply (*unfold HARoot-def*)
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold IsCompFun-def hierauto-def HierAuto-def*)
apply *fast*
done

lemma *HARoot-ran-CompFun2:*
 $S \in ran (CompFun HA) \longrightarrow HARoot HA \notin S$
apply (*unfold HARoot-def*)
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold IsCompFun-def hierauto-def HierAuto-def*)
apply *fast*
done

7.2.8 *CompFun*

lemma *IsCompFun-HA [simp]:*
 $IsCompFun (SAs HA) (CompFun HA)$
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold hierauto-def HierAuto-def*)
apply *auto*
done

lemma *dom-CompFun [simp]:*


```

  dom (CompFun HA) = HAStates HA
apply (cut-tac HA=HA in IsCompFun-HA)
apply (unfold IsCompFun-def HAStates-def)
apply auto
done

```

```

lemma ran-CompFun [simp]:
  Union (ran (CompFun HA)) = ((SAs HA) - {Root (SAs HA)(CompFun HA)})
apply (cut-tac HA=HA in IsCompFun-HA)
apply (unfold IsCompFun-def)
apply fast
done

```

```

lemma ran-CompFun-subseteq:
  Union (ran (CompFun HA))  $\subseteq$  (SAs HA)
apply (cut-tac HA=HA in IsCompFun-HA)
apply (unfold IsCompFun-def)
apply fast
done

```

```

lemma ran-CompFun-is-not-SA:
   $\neg$  Sas  $\subseteq$  (SAs HA)  $\implies$  Sas  $\notin$  (ran (CompFun HA))
apply (cut-tac HA=HA in IsCompFun-HA)
apply (unfold IsCompFun-def)
apply fast
done

```

```

lemma HAStates-HARoot-CompFun [simp]:
  S  $\in$  HAStates HA  $\implies$  HARoot HA  $\notin$  the (CompFun HA S)
apply (rule ran-dom-the)
back
apply (simp add: HARoot-ran-CompFun2 HARoot-def HAStates-def)+
done

```

```

lemma HAStates-CompFun-SAs:
  S  $\in$  HAStates A  $\implies$  the (CompFun A S)  $\subseteq$  SAs A
apply auto
apply (rename-tac T)
apply (cut-tac HA=A in ran-CompFun)
apply (erule equalityE)
apply (erule-tac c=T in subsetCE)
apply (drule ran-dom-the)
apply auto
done

```

```

lemma HAStates-CompFun-notmem [simp]:
   $\llbracket$  S  $\in$  HAStates A; SA  $\in$  the (CompFun A S)  $\rrbracket \implies$  S  $\notin$  States SA
apply (unfold HAStates-def)
apply auto

```

```

apply (rename-tac T)
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=SA in ballE)
apply (erule-tac x=T in ballE)
apply auto
prefer 2
apply (cut-tac A=A and S=S in HAStates-CompFun-SAs)
apply (unfold HAStates-def)
apply simp
apply fast
apply fast
apply (cut-tac HA=A in NoCycles-HA)
apply (unfold NoCycles-def)
apply (erule-tac x={S} in ballE)
apply auto
done

```

lemma *CompFun-Int-disjoint:*

$\llbracket S \neq T; S \in HAStates\ A; T \in HAStates\ A \rrbracket \implies \text{the } (CompFun\ A\ T) \cap \text{the } (CompFun\ A\ S) = \{\}$

```

apply auto
apply (rename-tac U)
apply (cut-tac HA=A in OneAncestor-HA)
apply (unfold OneAncestor-def)
apply (erule-tac x=U in ballE)
prefer 2
apply simp
apply (fold HARoot-def)
apply (frule HAStates-HARoot-CompFun)
apply simp
apply (frule HAStates-CompFun-SAs)
apply auto
apply (erule-tac x=S in allE)
apply (erule-tac x=T in allE)
apply auto
apply (cut-tac HA=A in NoCycles-HA)
apply (unfold NoCycles-def)
apply (simp only: HAStates-def)
apply safe
apply (erule-tac x={S} in ballE)
apply simp
apply fast
apply simp
apply (cut-tac HA=A in NoCycles-HA)
apply (unfold NoCycles-def)
apply (simp only: HAStates-def)
apply safe
apply (erule-tac x={T} in ballE)

```

apply *simp*
apply *fast*
apply *simp*
done

7.2.9 *SAs*

lemma *finite-SAs* [*simp*]:
finite (SAs HA)
apply (*cut-tac Rep-hierauto-select*)
apply (*unfold hierauto-def HierAuto-def*)
apply *fast*
done

lemma *HASates-SAs-disjunct*:
 $HASates HA1 \cap HASates HA2 = \{\} \implies SAs HA1 \cap SAs HA2 = \{\}$
apply (*unfold UNION-eq HASates-def Int-def*)
apply *auto*
apply (*rename-tac SA*)
apply (*cut-tac SA=SA in EX-State-SA*)
apply (*erule exE*)
apply *auto*
done

lemma *HASates-CompFun-SAs-mem* [*simp*]:
 $\llbracket S \in HASates A; T \in the (CompFun A S) \rrbracket \implies T \in SAs A$
apply (*cut-tac A=A and S=S in HASates-CompFun-SAs*)
apply *auto*
done

lemma *SAs-States-HASates*:
 $SA \in SAs A \implies States SA \subseteq HASates A$
by (*unfold HASates-def, auto*)

7.2.10 *HAINitState*

lemma *HAINitState-HARoot* [*simp*]:
 $HAINitState A \in States (HARoot A)$
by (*unfold HAINitState-def, auto*)

lemma *HAINitState-HARoot2* [*simp*]:
 $HAINitState A \in States (Root (SAs A) (CompFun A))$
by (*fold HARoot-def, simp*)

lemma *HAINitStates-HASates* [*simp*]:
 $HAINitStates A \subseteq HASates A$
apply (*unfold HAINitStates-def HASates-def*)
apply *auto*
done

lemma *HAINitStates-HAStates2* [simp]:
 $S \in \text{HAINitStates } A \implies S \in \text{HAStates } A$
apply (*cut-tac A=A in HAINitStates-HAStates*)
apply *fast*
done

lemma *HAINitState-HAStates* [simp]:
 $\text{HAINitState } A \in \text{HAStates } A$
apply (*unfold HAStates-def*)
apply *auto*
apply (*rule-tac x=HARoot A in beXI*)
apply *auto*
done

lemma *HAINitState-HAINitStates* [simp]:
 $\text{HAINitState } A \in \text{HAINitStates } A$
by (*unfold HAINitStates-def HAINitState-def, auto*)

lemma *CompFun-HAINitStates-HAStates* [simp]:
 $\llbracket S \in \text{HAStates } A; SA \in \text{the } (\text{CompFun } A \ S) \rrbracket \implies (\text{InitState } SA) \in \text{HAINitStates } A$
apply (*unfold HAINitStates-def*)
apply *auto*
done

lemma *CompFun-HAINitState-HAINitStates* [simp]:
 $\llbracket SA \in \text{the } (\text{CompFun } A \ (\text{HAINitState } A)) \rrbracket \implies (\text{InitState } SA) \in \text{HAINitStates } A$
apply (*unfold HAINitStates-def*)
apply *auto*
apply (*rule-tac x=SA in beXI*)
apply *auto*
apply (*cut-tac A=A and S=HAINitState A in HAStates-CompFun-SAs*)
apply *auto*
done

lemma *HAINitState-notmem-States* [simp]:
 $\llbracket S \in \text{HAStates } A; SA \in \text{the } (\text{CompFun } A \ S) \rrbracket \implies \text{HAINitState } A \notin \text{States } SA$
apply (*cut-tac HA=A in MutuallyDistinct-HA*)
apply (*unfold MutuallyDistinct-def*)
apply (*erule-tac x=SA in ballE*)
apply (*erule-tac x=HARoot A in ballE*)
apply *auto*
done

lemma *InitState-notmem-States* [simp]:
 $\llbracket S \in \text{HAStates } A; SA \in \text{the } (\text{CompFun } A \ S);$
 $T \in \text{HAINitStates } A; T \neq \text{InitState } SA \rrbracket$
 $\implies T \notin \text{States } SA$

```

apply (unfold HAINitStates-def)
apply auto
apply (rename-tac SAA)
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=SA in ballE)
apply (erule-tac x=SAA in ballE)
apply auto
done

```

```

lemma InitState-States-notmem [simp]:
   $\llbracket B \in \text{SAs } A; C \in \text{SAs } A; B \neq C \rrbracket \implies \text{InitState } B \notin \text{States } C$ 
apply auto
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply force
done

```

```

lemma OneHAINitState-SASates:
   $\llbracket S \in \text{HAINitStates } A; T \in \text{HAINitStates } A; \\ S \in \text{States } SA; T \in \text{States } SA; SA \in \text{SAs } A \rrbracket \implies \\ S = T$ 
apply (unfold HAINitStates-def)
apply auto
apply (rename-tac AA AAA)
apply (case-tac AA = SA)
apply auto
apply (case-tac AAA = SA)
apply auto
done

```

7.2.11 Chi

```

lemma HARootStates-notmem-Chi [simp]:
   $\llbracket S \in \text{HASates } A; T \in \text{States } (\text{HARoot } A) \rrbracket \implies T \notin \text{Chi } A S$ 
apply (unfold Chi-def restrict-def, auto)
apply (rename-tac SA)
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=HARoot A in ballE)
apply (erule-tac x=SA in ballE)
apply auto
done

```

```

lemma SASates-notmem-Chi [simp]:
   $\llbracket S \in \text{States } SA; T \in \text{States } SA; \\ SA \in \text{SAs } A \rrbracket \implies T \notin \text{Chi } A S$ 
apply (unfold Chi-def restrict-def, auto)
apply (rename-tac SAA)

```

apply (*cut-tac HA=A in MutuallyDistinct-HA*)
apply (*unfold MutuallyDistinct-def*)
apply (*erule-tac x=SAA in ballE*)
apply (*erule-tac x=SA in ballE*)
apply *auto*
apply (*unfold HAStates-def*)
apply *auto*
done

lemma *HAInitState-notmem-Chi* [*simp*]:
 $S \in HAStates\ A \implies HAINitState\ A \notin Chi\ A\ S$
by (*unfold Chi-def restrict-def, auto*)

lemma *Chi-HAStates* [*simp*]:
 $T \in HAStates\ A \implies (Chi\ A\ T) \subseteq HAStates\ A$
apply (*unfold Chi-def restrict-def*)
apply (*auto*)
apply (*cut-tac A=A and S=T in HAStates-CompFun-SAs*)
apply (*unfold HAStates-def*)
apply *auto*
done

lemma *Chi-HAStates-Self* [*simp*]:
 $s \in HAStates\ a \implies s \notin (Chi\ a\ s)$
by (*unfold Chi-def restrict-def, auto*)

lemma *ChiRel-HAStates-Self* [*simp*]:
 $(s,s) \notin (ChiRel\ a)$
by (*unfold ChiRel-def, auto*)

lemma *HAStates-Chi-NoCycles*:
 $\llbracket s \in HAStates\ a; t \in HAStates\ a; s \in Chi\ a\ t \rrbracket \implies t \notin Chi\ a\ s$
apply (*unfold Chi-def restrict-def*)
apply *auto*
apply (*cut-tac HA=a in NoCycles-HA*)
apply (*unfold NoCycles-def*)
apply (*erule-tac x={s,t} in ballE*)
apply *auto*
done

lemma *HAStates-Chi-NoCycles-trans*:
 $\llbracket s \in HAStates\ a; t \in HAStates\ a; u \in HAStates\ a;$
 $t \in Chi\ a\ s; u \in Chi\ a\ t \rrbracket \implies s \notin Chi\ a\ u$
apply (*unfold Chi-def restrict-def*)
apply *auto*
apply (*cut-tac HA=a in NoCycles-HA*)
apply (*unfold NoCycles-def*)
apply (*erule-tac x={s,t,u} in ballE*)
prefer 2

```

apply simp
apply (unfold HAStates-def)
apply auto
done

```

```

lemma Chi-range-disjoint:
   $\llbracket S \neq T; T \in HAStates\ A; S \in HAStates\ A; U \in Chi\ A\ S \rrbracket \implies U \notin Chi\ A\ T$ 
apply (frule CompFun-Int-disjoint)
apply auto
apply (unfold Chi-def restrict-def)
apply auto
apply (rename-tac SA SAA)
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=SA in ballE)
apply (erule-tac x=SAA in ballE)
apply auto
done

```

```

lemma SASates-Chi-trans [rule-format]:
   $\llbracket U \in Chi\ A\ T; S \in Chi\ A\ U; T \in States\ SA; SA \in SAs\ A; U \in HAStates\ A \rrbracket \implies S \notin States\ SA$ 
apply (frule HAStates-SA-mem)
apply auto
apply (unfold Chi-def restrict-def)
apply auto
apply (rename-tac SAA SAAA)
apply (cut-tac HA=A in NoCycles-HA)
apply (unfold NoCycles-def)
apply (erule-tac x={U,T} in ballE)
prefer 2
apply (simp only: HAStates-def)
apply auto
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (rotate-tac -1)
apply (erule-tac x=SA in ballE)
apply (rotate-tac -1)
apply (erule-tac x=SAAA in ballE)
apply auto
done

```

7.2.12 *ChiRel*

```

lemma finite-ChiRel [simp]:
  finite (ChiRel A)
apply (rule-tac B=HAStates A  $\times$  HAStates A in finite-subset)
apply auto
done

```

lemma *ChiRel-HAStates-subseteq* [simp]:
 $(\text{ChiRel } A) \subseteq (\text{HAStates } A \times \text{HAStates } A)$
apply (unfold ChiRel-def Chi-def restrict-def)
apply auto
done

lemma *ChiRel-CompFun*:
 $s \in \text{HAStates } a \implies \text{ChiRel } a \text{ “ } \{s\} = (\bigcup x \in \text{the } (\text{CompFun } a \ s)). \text{ States } x$
apply (unfold ChiRel-def Chi-def restrict-def Image-def)
apply simp
apply auto
apply (frule HAStates-CompFun-SAs-mem)
apply fast
apply (unfold HAStates-def)
apply fast
done

lemma *ChiRel-HARoot*:
 $\llbracket (x,y) \in \text{ChiRel } A \rrbracket \implies y \notin \text{States } (\text{HARoot } A)$
apply (unfold ChiRel-def Chi-def)
apply auto
apply (unfold restrict-def)
apply auto
apply (rename-tac SA)
apply (frule HAStates-HARoot-CompFun)
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply auto
apply (erule-tac x=SA in ballE)
apply (erule-tac x=HARoot A in ballE)
apply auto
done

lemma *HAStates-CompFun-States-ChiRel*:
 $S \in \text{HAStates } A \implies \bigcup (\text{States } \text{“ the } (\text{CompFun } A \ S)) = \text{ChiRel } A \text{ “ } \{S\}$
apply (unfold ChiRel-def Chi-def restrict-def)
apply auto
apply (drule HAStates-CompFun-SAs)
apply (subst HAStates-def)
apply fast
done

lemma *HAInitState-notmem-Range-ChiRel* [simp]:
 $\text{HAInitState } A \notin \text{Range } (\text{ChiRel } A)$
by (unfold ChiRel-def, auto)

lemma *HAInitState-notmem-Range-ChiRel2* [simp]:
 $(S, \text{HAInitState } A) \notin (\text{ChiRel } A)$

by (unfold ChiRel-def, auto)

lemma *ChiRel-OneAncestor-notmem:*

$\llbracket S \neq T; (S,U) \in \text{ChiRel } A \rrbracket \implies (T,U) \notin \text{ChiRel } A$
apply (unfold ChiRel-def)
apply auto
apply (simp only: Chi-range-disjoint)
done

lemma *ChiRel-OneAncestor:*

$\llbracket (S1,U) \in \text{ChiRel } A; (S2,U) \in \text{ChiRel } A \rrbracket \implies S1 = S2$
apply (rule notnotD, rule notI)
apply (simp add: ChiRel-OneAncestor-notmem)
done

lemma *CompFun-ChiRel:*

$\llbracket S1 \in \text{HAStates } A; SA \in \text{the } (\text{CompFun } A \ S1);$
 $S2 \in \text{States } SA \rrbracket \implies (S1,S2) \in \text{ChiRel } A$
apply (unfold ChiRel-def Chi-def restrict-def)
apply auto
apply (cut-tac A=A and S=S1 in HAStates-CompFun-SAs)
apply (unfold HAStates-def)
apply auto
done

lemma *CompFun-ChiRel2:*

$\llbracket (S,T) \in \text{ChiRel } A; T \in \text{States } SA; SA \in \text{SAs } A \rrbracket \implies SA \in \text{the } (\text{CompFun } A \ S)$
apply (unfold ChiRel-def Chi-def restrict-def)
apply auto
apply (rename-tac SAA)
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=SA in ballE)
apply (rotate-tac -1)
apply (erule-tac x=SAA in ballE)
apply auto
done

lemma *ChiRel-HAStates-NoCycles:*

$(s,t) \in (\text{ChiRel } a) \implies (t,s) \notin (\text{ChiRel } a)$
apply (unfold ChiRel-def)
apply auto
apply (frule HAStates-Chi-NoCycles)
apply auto
done

lemma *HAStates-ChiRel-NoCycles-trans:*

$\llbracket (s,t) \in (\text{ChiRel } a); (t,u) \in (\text{ChiRel } a) \rrbracket \implies (u,s) \notin (\text{ChiRel } a)$
apply (unfold ChiRel-def)

apply *auto*
apply (*frule* *HAStates-Chi-NoCycles-trans*)
apply *fast*
back
back
prefer 3
apply *fast*
apply *auto*
done

lemma *SAStates-ChiRel*:
 $\llbracket S \in \text{States } SA; T \in \text{States } SA;$
 $SA \in \text{SAs } A \rrbracket \implies (S, T) \notin (\text{ChiRel } A)$
by (*unfold ChiRel-def, auto*)

lemma *ChiRel-SA-OneAncestor*:
 $\llbracket (S, T) \in \text{ChiRel } A; T \in \text{States } SA;$
 $U \in \text{States } SA; SA \in \text{SAs } A \rrbracket \implies$
 $(S, U) \in \text{ChiRel } A$
apply (*frule CompFun-ChiRel2*)
apply *auto*
apply (*rule CompFun-ChiRel*)
apply *auto*
done

lemma *ChiRel-OneAncestor2*:
 $\llbracket S \in \text{HAStates } A; S \notin \text{States } (\text{HARoot } A) \rrbracket \implies$
 $\exists! T. (T, S) \in \text{ChiRel } A$
apply (*unfold ChiRel-def*)
apply *auto*
prefer 2
apply (*rename-tac T U*)
prefer 2
apply (*unfold Chi-def restrict-def*)
apply *auto*
prefer 2
apply (*rename-tac SA SAA*)
prefer 2
apply (*cut-tac HA=A in OneAncestor-HA*)
apply (*unfold OneAncestor-def*)
apply (*fold HARoot-def*)
apply *auto*
apply (*simp cong: rev-conj-cong*)
apply (*unfold HAStates-def*)
apply *auto*
apply (*rename-tac SA*)
apply (*erule-tac x=SA in ballE*)
apply *auto*
apply (*case-tac T = U*)

```

apply auto
apply (frule CompFun-Int-disjoint)
apply (unfold HASStates-def)
apply auto
apply (case-tac SA=SAA)
apply auto
apply (cut-tac HA=A in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=SAA in ballE)
apply (erule-tac x=SA in ballE)
apply auto
apply (cut-tac S=T and A=A in HASStates-CompFun-SAs)
apply (unfold HASStates-def)
apply fast
apply fast
apply (cut-tac S=U and A=A in HASStates-CompFun-SAs)
apply (unfold HASStates-def)
apply fast
apply fast
done

```

lemma *HARootStates-notmem-Range-ChiRel* [*simp*]:
 $S \in \text{States } (\text{HARoot } A) \implies S \notin \text{Range } (\text{ChiRel } A)$
by (*unfold* *ChiRel-def*, *auto*)

lemma *ChiRel-int-disjoint*:
 $S \neq T \implies (\text{ChiRel } A \text{ `` } \{S\}) \cap (\text{ChiRel } A \text{ `` } \{T\}) = \{\}$
apply (*unfold* *ChiRel-def*)
apply *auto*
apply (*simp only: Chi-range-disjoint*)
done

lemma *SASStates-ChiRel-trans* [*rule-format*]:
 $\llbracket (S,U) \in (\text{ChiRel } A); (U,T) \in \text{ChiRel } A; S \in \text{States } SA; SA \in \text{SAs } A \rrbracket \implies T \notin \text{States } SA$
apply *auto*
apply (*unfold* *ChiRel-def*)
apply *auto*
apply (*frule* *SASStates-Chi-trans*)
back
apply *fast+*
done

lemma *HAINitStates-InitState-trancl*:
 $\llbracket S \in \text{HAINitStates } (\text{HA } ST); A \in \text{the } (\text{CompFun } (\text{HA } ST) S) \rrbracket \implies$
 $(S, \text{InitState } A) \in (\text{ChiRel } (\text{HA } ST) \cap \text{HAINitStates } (\text{HA } ST) \times \text{HAINitStates } (\text{HA } ST))^+$
apply (*case-tac* $S \in \text{HASStates } (\text{HA } ST)$)
apply (*frule* *CompFun-ChiRel*)

```

apply fast+
apply (rule InitState-States)
apply auto
apply (rule r-into-trancl')
apply auto
apply (rule CompFun-HAInitStates-HAStates)
apply auto
done

```

lemma *HAInitStates-InitState-trancl2*:

```

 $\llbracket S \in HAStates (HA ST); A \in the (CompFun (HA ST) S);$ 
 $(x, S) \in (ChiRel (HA ST) \cap HAINitStates (HA ST) \times HAINitStates (HA ST))^+$ 
 $\rrbracket$ 
 $\implies (x, InitState A) \in (ChiRel (HA ST) \cap HAINitStates (HA ST) \times HAINitStates$ 
 $(HA ST))^+$ 
apply (rule-tac a=x and b=S and r=ChiRel (HA ST) \cap HAINitStates (HA ST)
 $\times HAINitStates (HA ST)$  in converse-trancl-induct)
apply auto
prefer 2
apply (rename-tac T U)
prefer 2
apply (case-tac S \in HAStates (HA ST))
apply (frule CompFun-ChiRel)
apply fast
apply (rule InitState-States)
apply simp
apply (rule trancl-trans [of - S])
apply (rule r-into-trancl')
apply auto
apply (rule r-into-trancl')
apply auto
apply (rule CompFun-HAInitStates-HAStates)
prefer 2
apply fast
apply (cut-tac A=HA ST in HAINitStates-HAStates, fast)
apply (rule-tac y = U in trancl-trans)
apply (rule r-into-trancl')
apply auto
done

```

7.2.13 *ChiPlus*

lemma *ChiPlus-ChiRel [simp]*:

```

 $(S, T) \in ChiRel A \implies (S, T) \in ChiPlus A$ 
apply (unfold ChiPlus-def)
apply (frule r-into-trancl)
apply auto
done

```

lemma *ChiPlus-HAStates* [simp]:
 $(\text{ChiPlus } A) \subseteq (\text{HAStates } A \times \text{HAStates } A)$
apply (unfold *ChiPlus-def*)
apply (rule *trancl-subset-Sigma*)
apply auto
done

lemma *ChiPlus-subset-States*:
 $\text{ChiPlus } a \subseteq \bigcup (\text{States } (SAs\ a))$
apply (cut-tac $A=a$ in *ChiPlus-HAStates*)
apply (unfold *HAStates-def*)
apply auto
done

lemma *finite-ChiPlus* [simp]:
 $\text{finite } (\text{ChiPlus } A)$
apply (rule-tac $B=\text{HAStates } A \times \text{HAStates } A$ in *finite-subset*)
apply auto
done

lemma *ChiPlus-OneAncestor*:
 $\llbracket S \in \text{HAStates } A; S \notin \text{States } (\text{HARoot } A) \rrbracket \implies$
 $\exists T. (T, S) \in \text{ChiPlus } A$
apply (unfold *ChiPlus-def*)
apply (rule *ChiRel-OneAncestor2*)
apply auto
done

lemma *ChiPlus-HAStates-Left*:
 $(S, T) \in \text{ChiPlus } A \implies S \in \text{HAStates } A$
apply (cut-tac $A=A$ in *ChiPlus-HAStates*)
apply (unfold *HAStates-def*)
apply auto
done

lemma *ChiPlus-HAStates-Right*:
 $(S, T) \in \text{ChiPlus } A \implies T \in \text{HAStates } A$
apply (cut-tac $A=A$ in *ChiPlus-HAStates*)
apply (unfold *HAStates-def*)
apply auto
done

lemma *ChiPlus-ChiRel-int* [rule-format]:
 $\llbracket (T, S) \in (\text{ChiPlus } A) \rrbracket \implies (\text{ChiPlus } A \subseteq \{T\}) \cap (\text{ChiRel } A \subseteq \{S\}) = (\text{ChiRel } A \subseteq \{S\})$
apply (unfold *ChiPlus-def*)
apply (rule-tac $a=T$ and $b=S$ and $r=(\text{ChiRel } A)$ in *converse-trancl-induct*)
apply auto
done

lemma *ChiPlus-ChiPlus-int* [rule-format]:
 $\llbracket (T,S) \in (\text{ChiPlus } A) \rrbracket \implies (\text{ChiPlus } A \text{ `` } \{T\}) \cap (\text{ChiPlus } A \text{ `` } \{S\}) = (\text{ChiPlus } A \text{ `` } \{S\})$
apply (*unfold ChiPlus-def*)
apply (*rule-tac a=T and b=S and r=(ChiRel A) in converse-trancl-induct*)
apply *auto*
done

lemma *ChiPlus-ChiRel-NoCycle-1* [rule-format]:
 $\llbracket (T,S) \in \text{ChiPlus } A \rrbracket \implies$
 $(\text{insert } S (\text{insert } T (\{U. (T,U) \in \text{ChiPlus } A \wedge (U,S) \in \text{ChiPlus } A\}))) \cap (\text{ChiRel } A \text{ `` } \{T\}) \neq \{\}$
apply (*unfold ChiPlus-def*)
apply (*rule-tac a=T and b=S and r=(ChiRel A) in converse-trancl-induct*)
apply (*unfold Image-def Int-def*)
apply *auto*
done

lemma *ChiPlus-ChiRel-NoCycle-2* [rule-format]:
 $\llbracket (T,S) \in \text{ChiPlus } A \rrbracket \implies (S,T) \in (\text{ChiRel } A) \longrightarrow$
 $(\text{insert } S (\text{insert } T (\{U. (T,U) \in \text{ChiPlus } A \wedge (U,S) \in \text{ChiPlus } A\}))) \cap (\text{ChiRel } A \text{ `` } \{S\}) \neq \{\}$
apply (*unfold ChiPlus-def*)
apply (*rule-tac a=T and b=S and r=(ChiRel A) in converse-trancl-induct*)
apply (*unfold Image-def Int-def*)
apply *auto*
done

lemma *ChiPlus-ChiRel-NoCycle-3* [rule-format]:
 $\llbracket (T,S) \in \text{ChiPlus } A \rrbracket \implies (S,T) \in (\text{ChiRel } A) \longrightarrow (T,U) \in \text{ChiPlus } A \longrightarrow (U,$
 $S) \in \text{ChiPlus } A \longrightarrow$
 $(\text{insert } S (\text{insert } T (\{U. (T,U) \in \text{ChiPlus } A \wedge (U,S) \in \text{ChiPlus } A\}))) \cap (\text{ChiRel } A \text{ `` } \{U\}) \neq \{\}$
apply (*unfold ChiPlus-def*)
apply (*rule-tac a=T and b=S and r=(ChiRel A) in trancl-induct*)
apply (*unfold Image-def Int-def, simp*)
apply (*rename-tac V*)
prefer 2
apply (*rename-tac V W*)
prefer 2
apply (*simp, safe*)
apply (*simp only: ChiRel-HAStates-NoCycles*)
apply *simp*
apply (*case-tac (U,W) \in (ChiRel A), fast, rotate-tac 5, frule tranclD3, fast, blast*
intro: trancl-into-trancl)
done

lemma *ChiPlus-ChiRel-NoCycle-4* [rule-format]:

```

[[ (T,S) ∈ ChiPlus A ] ⇒ (S,T) ∈ (ChiRel A) → ((ChiPlus A “{T}) ∩ (ChiRel
A “{S})) ≠ {}
apply (unfold ChiPlus-def)
apply (rule-tac a=T and b=S and r=(ChiRel A) in trancl-induct)
apply (unfold Image-def Int-def)
apply auto
apply (simp only: ChiRel-HAStates-NoCycles)
apply (rule-tac x=T in exI)
apply simp
apply (rule-tac x=T in exI)
apply simp
done

```

lemma *ChiRel-ChiPlus-NoCycles*:

```

(S,T) ∈ (ChiRel A) ⇒ (T,S) ∉ (ChiPlus A)
apply (cut-tac HA=A in NoCycles-HA)
apply (unfold NoCycles-def)
apply (erule-tac x=insert S (insert T ({U. (T,U) ∈ ChiPlus A ∧ (U,S) ∈ ChiPlus
A})) in ballE)
prefer 2
apply (simp add: ChiPlus-subset-States)
apply (cut-tac A=A in ChiPlus-HAStates)
apply (unfold HAStates-def)
apply auto
apply (frule ChiPlus-ChiRel-NoCycle-2)
apply fast
apply (simp add: ChiRel-CompFun)
apply (frule ChiPlus-ChiRel-NoCycle-1)
apply (simp add: ChiRel-CompFun)
apply (frule ChiPlus-ChiRel-NoCycle-3)
apply fast
apply fast
back
apply fast
apply (rename-tac V)
apply (case-tac V ∈ HAStates A)
apply (simp add: ChiRel-CompFun)
apply (simp only: ChiPlus-HAStates-Right)
apply fast
done

```

lemma *ChiPlus-ChiPlus-NoCycles*:

```

(S,T) ∈ (ChiPlus A) ⇒ (T,S) ∉ (ChiPlus A)
apply (unfold ChiPlus-def)
apply (rule-tac a=S and b=T and r=(ChiRel A) in trancl-induct)
apply fast
apply (frule ChiRel-ChiPlus-NoCycles)
apply (auto intro: trancl-into-trancl2 simp add: ChiPlus-def)
done

```

lemma *ChiPlus-NoCycles* [rule-format]:
 $(S,T) \in (\text{ChiPlus } A) \implies S \neq T$
apply (frule *ChiPlus-ChiPlus-NoCycles*)
apply auto
done

lemma *ChiPlus-NoCycles-2* [simp]:
 $(S,S) \notin (\text{ChiPlus } A)$
apply (rule notI)
apply (frule *ChiPlus-NoCycles*)
apply fast
done

lemma *ChiPlus-ChiPlus-NoCycles-2*:
 $\llbracket (S,U) \in \text{ChiPlus } A; (U,T) \in \text{ChiPlus } A \rrbracket \implies (T,S) \notin \text{ChiPlus } A$
apply (rule *ChiPlus-ChiPlus-NoCycles*)
apply (auto intro: trancl-trans simp add: *ChiPlus-def*)
done

lemma *ChiRel-ChiPlus-trans*:
 $\llbracket (U,S) \in \text{ChiPlus } A; (S,T) \in \text{ChiRel } A \rrbracket \implies (U,T) \in \text{ChiPlus } A$
apply (unfold *ChiPlus-def*)
apply auto
done

lemma *ChiRel-ChiPlus-trans2*:
 $\llbracket (U,S) \in \text{ChiRel } A; (S,T) \in \text{ChiPlus } A \rrbracket \implies (U,T) \in \text{ChiPlus } A$
apply (unfold *ChiPlus-def*)
apply auto
done

lemma *ChiPlus-ChiRel-Ex* [rule-format]:
 $\llbracket (S,T) \in \text{ChiPlus } A \rrbracket \implies (S,T) \notin \text{ChiRel } A \longrightarrow$
 $(\exists U. (S,U) \in \text{ChiPlus } A \wedge (U,T) \in \text{ChiRel } A)$
apply (unfold *ChiPlus-def*)
apply (rule-tac a=S and b=T and r=(*ChiRel A*) in converse-trancl-induct)
apply auto
apply (rename-tac U)
apply (rule-tac x=U in exI)
apply auto
done

lemma *ChiPlus-ChiRel-Ex2* [rule-format]:
 $\llbracket (S,T) \in \text{ChiPlus } A \rrbracket \implies (S,T) \notin \text{ChiRel } A \longrightarrow$
 $(\exists U. (S,U) \in \text{ChiRel } A \wedge (U,T) \in \text{ChiPlus } A)$
apply (unfold *ChiPlus-def*)
apply (rule-tac a=S and b=T and r=(*ChiRel A*) in converse-trancl-induct)
apply auto

done

lemma *HARootStates-Range-ChiPlus* [simp]:

$\llbracket S \in \text{States } (\text{HARoot } A) \rrbracket \implies S \notin \text{Range } (\text{ChiPlus } A)$
by (*unfold ChiPlus-def*, *auto*)

lemma *HARootStates-Range-ChiPlus2* [simp]:

$\llbracket S \in \text{States } (\text{HARoot } A) \rrbracket \implies (x,S) \notin (\text{ChiPlus } A)$
by (*frule HARootStates-Range-ChiPlus*, *unfold Domain-converse* [symmetric], *fast*)

lemma *SASates-ChiPlus-ChiRel-NoCycle-1* [rule-format]:

$\llbracket (S,U) \in \text{ChiPlus } A; SA \in \text{SAs } A \rrbracket \implies (U,T) \in (\text{ChiRel } A) \longrightarrow S \in \text{States } SA$
 $\longrightarrow T \in \text{States } SA \longrightarrow$

$(\text{insert } S (\text{insert } U (\{V. (S,V) \in \text{ChiPlus } A \wedge (V,U) \in \text{ChiPlus } A\}))) \cap (\text{ChiRel } A \text{ “ } \{U\}) \neq \{\}$

apply (*unfold ChiPlus-def*)

apply (*rule-tac a=S and b=U and r=(ChiRel A) in converse-trancl-induct*)

apply (*simp*, *safe*)

apply (*simp only: SASates-ChiRel-trans*)

apply (*simp add: ChiRel-CompFun*)

apply *safe*

apply (*erule-tac x=SA in ballE*)

apply (*simp add: CompFun-ChiRel2*) $+$

apply (*simp add: Int-def*, *fast*)

apply *auto*

apply (*fold ChiPlus-def*)

apply (*rename-tac W*)

apply (*frule-tac U=U and T=U and S=W in ChiRel-ChiPlus-trans2*)

apply *auto*

done

lemma *SASates-ChiPlus-ChiRel-NoCycle-2* [rule-format]:

$\llbracket (S,U) \in \text{ChiPlus } A \rrbracket \implies (U,T) \in (\text{ChiRel } A) \longrightarrow$

$(\text{insert } S (\text{insert } U (\{V. (S,V) \in \text{ChiPlus } A \wedge (V,U) \in \text{ChiPlus } A\}))) \cap (\text{ChiRel } A \text{ “ } \{S\}) \neq \{\}$

apply (*unfold ChiPlus-def*)

apply (*rule-tac a=S and b=U and r=(ChiRel A) in converse-trancl-induct*)

apply (*unfold Image-def Int-def*)

apply *auto*

done

lemma *SASates-ChiPlus-ChiRel-NoCycle-3* [rule-format]:

$\llbracket (S,U) \in \text{ChiPlus } A \rrbracket \implies (U,T) \in (\text{ChiRel } A) \longrightarrow (S,s) \in \text{ChiPlus } A \longrightarrow$
 $(s,U) \in \text{ChiPlus } A \longrightarrow$

$(\text{insert } S (\text{insert } U (\{V. (S,V) \in \text{ChiPlus } A \wedge (V,U) \in \text{ChiPlus } A\}))) \cap (\text{ChiRel } A \text{ “ } \{s\}) \neq \{\}$

apply (*unfold ChiPlus-def*)

```

apply (rule-tac  $a=S$  and  $b=U$  and  $r=(ChiRel\ A)$  in trancl-induct)
apply fast
apply (rename-tac  $W$ )
prefer 2
apply (rename-tac  $W\ X$ )
prefer 2
apply (unfold Image-def Int-def)
apply (simp, safe)
apply (fold ChiPlus-def)
apply (case-tac  $(s, W) \in ChiRel\ A$ )
apply fast
apply (frule-tac  $S=s$  and  $T=W$  in ChiPlus-ChiRel-Ex2)
apply simp
apply safe
apply (rename-tac  $X$ )
apply (rule-tac  $x=X$  in exI)
apply (fast intro: ChiRel-ChiPlus-trans)
apply simp
apply (case-tac  $(s, X) \in ChiRel\ A$ )
apply force
apply (frule-tac  $S=s$  and  $T=X$  in ChiPlus-ChiRel-Ex2)
apply simp
apply safe
apply (rename-tac  $Y$ )
apply (erule-tac  $x=Y$  in allE)
apply simp
apply (fast intro: ChiRel-ChiPlus-trans)
apply simp
apply (case-tac  $(s, X) \in ChiRel\ A$ )
apply force
apply (frule-tac  $S=s$  and  $T=X$  in ChiPlus-ChiRel-Ex2)
apply simp
apply safe
apply (rename-tac  $Y$ )
apply (erule-tac  $x=Y$  in allE)
apply simp
apply (fast intro: ChiRel-ChiPlus-trans)
apply fastforce
apply simp
apply (erule-tac  $x=W$  in allE)
apply simp
apply simp
apply (rename-tac  $Y$ )
apply (erule-tac  $x=Y$  in allE)
apply simp
apply (fast intro: ChiRel-ChiPlus-trans)
done

```

lemma *SASates-ChiPlus-ChiRel-trans* [*rule-format*]:

```

[[ (S,U) ∈ (ChiPlus A); (U,T) ∈ (ChiRel A); S ∈ States SA; SA ∈ SAs A ]] ⇒
T ∉ States SA
apply (cut-tac HA=A in NoCycles-HA)
apply (unfold NoCycles-def)
apply (erule-tac x=insert S (insert U ({ V. (S,V) ∈ ChiPlus A ∧ (V,U) ∈ ChiPlus
A})) in ballE)
prefer 2
apply (simp add: ChiPlus-subset-States)
apply (cut-tac A=A in ChiPlus-HAStates)
apply (unfold HAStates-def)
apply auto[1]
apply safe
apply fast
apply (frule SAStates-ChiPlus-ChiRel-NoCycle-2)
apply fast
apply (frule HAStates-SA-mem)
apply fast
apply (simp only: ChiRel-CompFun)
apply (frule SAStates-ChiPlus-ChiRel-NoCycle-1)
apply auto[3]
apply fast
apply (simp add: ChiRel-CompFun)
apply (frule SAStates-ChiPlus-ChiRel-NoCycle-3)
apply fast
apply fast
back
apply fast
apply (simp only: ChiPlus-HAStates-Left ChiRel-CompFun)
done

```

```

lemma SAStates-ChiPlus2 [rule-format]:
[[ (S,T) ∈ ChiPlus A; SA ∈ SAs A ]] ⇒ S ∈ States SA → T ∉ States SA
apply (unfold ChiPlus-def)
apply (rule-tac a=S and b=T and r=(ChiRel A) in trancl-induct)
apply auto
apply (rename-tac U)
apply (frule-tac S=S and T=U in SAStates-ChiRel)
apply auto
apply (fold ChiPlus-def)
apply (simp only: SAStates-ChiPlus-ChiRel-trans)
done

```

```

lemma SAStates-ChiPlus [rule-format]:
[[ S ∈ States SA; T ∈ States SA; SA ∈ SAs A ]] ⇒ (S,T) ∉ ChiPlus A
apply auto
apply (simp only: SAStates-ChiPlus2)
done

```

```

lemma SAStates-ChiPlus-ChiRel-OneAncestor [rule-format]:

```

```

  [| T ∈ States SA; SA ∈ SAs A; (S,U) ∈ ChiPlus A ] ==> S ≠ T → S ∈ States
  SA → (T,U) ∉ ChiRel A
  apply (unfold ChiPlus-def)
  apply (rule-tac a=S and b=U and r=(ChiRel A) in trancl-induct)
  apply auto
  apply (simp add: ChiRel-OneAncestor-notmem)
  apply (rename-tac V W)
  apply (fold ChiPlus-def)
  apply (case-tac V=T)
  apply (simp add: ChiRel-OneAncestor-notmem SASates-ChiPlus)+
  done

```

```

lemma SASates-ChiPlus-OneAncestor [rule-format]:
  [| T ∈ States SA; SA ∈ SAs A; (S,U) ∈ ChiPlus A ] ==> S ≠ T →
  S ∈ States SA → (T,U) ∉ ChiPlus A
  apply (unfold ChiPlus-def)
  apply (rule-tac a=S and b=U and r=(ChiRel A) in trancl-induct)
  apply auto
  apply (fold ChiPlus-def)
  apply (rename-tac V)
  apply (frule-tac T=S and S=T and U=V in SASates-ChiPlus-ChiRel-OneAncestor)
  apply auto
  apply (rename-tac V W)
  apply (frule-tac S=T and T=W in ChiPlus-ChiRel-Ex)
  apply auto
  apply (frule-tac T=T and S=S and U=W in SASates-ChiPlus-ChiRel-OneAncestor)
  apply auto
  apply (rule ChiRel-ChiPlus-trans)
  apply auto
  apply (rename-tac X)
  apply (case-tac V=X)
  apply simp
  apply (simp add: ChiRel-OneAncestor-notmem)
  done

```

```

lemma ChiRel-ChiPlus-OneAncestor [rule-format]:
  [| (T,U) ∈ ChiPlus A ] ==> T ≠ S → (S,U) ∈ ChiRel A → (T,S) ∈ ChiPlus
  A
  apply (unfold ChiPlus-def)
  apply (rule-tac a=T and b=U and r=(ChiRel A) in trancl-induct)
  apply auto
  apply (fast intro: ChiRel-OneAncestor)
  apply (rename-tac V W)
  apply (case-tac S=V)
  apply auto
  apply (fast intro: ChiRel-OneAncestor)
  done

```

```

lemma ChiPlus-SA-OneAncestor [rule-format]:

```

```

    [[ (S,T) ∈ ChiPlus A;
      U ∈ States SA; SA ∈ SAs A ]] ⇒ T ∈ States SA →
      (S,U) ∈ ChiPlus A
apply (unfold ChiPlus-def)
apply (rule-tac a=S and b=T and r=(ChiRel A) in converse-trancl-induct)
apply auto
apply (frule ChiRel-SA-OneAncestor)
apply fast+
done

```

7.2.14 ChiStar

```

lemma ChiPlus-ChiStar [simp]:
  [[ (S,T) ∈ ChiPlus A ]] ⇒ (S,T) ∈ ChiStar A
by (unfold ChiPlus-def ChiStar-def, auto)

```

```

lemma HARootState-Range-ChiStar [simp]:
  [[ x ≠ S; S ∈ States (HARoot A) ]] ⇒ (x,S) ∉ (ChiStar A)
apply (unfold ChiStar-def)
apply (subst rtrancl-eq-or-trancl)
apply (fold ChiPlus-def)
apply auto
done

```

```

lemma ChiStar-Self [simp]:
  (S,S) ∈ ChiStar A
apply (unfold ChiStar-def)
apply simp
done

```

```

lemma ChiStar-Image [simp]:
  S ∈ M ⇒ S ∈ (ChiStar A “ M)
apply (unfold Image-def)
apply (auto intro: ChiStar-Self)
done

```

```

lemma ChiStar-ChiPlus-noteq:
  [[ S ≠ T; (S,T) ∈ ChiStar A ]] ⇒ (S,T) ∈ ChiPlus A
apply (unfold ChiPlus-def ChiStar-def)
apply (simp add: rtrancl-eq-or-trancl)
done

```

```

lemma ChiRel-ChiStar-trans:
  [[ (S,U) ∈ ChiStar A; (U,T) ∈ ChiRel A ]] ⇒ (S,T) ∈ ChiStar A
apply (unfold ChiStar-def)
apply auto
done

```

7.2.15 *InitConf*

```
lemma InitConf-HAStates [simp]:  
  InitConf A  $\subseteq$  HAStates A  
apply (unfold InitConf-def HAStates-def)  
apply auto  
apply (rule rtrancl-induct)  
back  
apply auto  
apply (rule-tac x=HARoot A in bexI)  
apply auto  
apply (unfold HAStates-def ChiRel-def)  
apply auto  
done
```

```
lemma InitConf-HAStates2 [simp]:  
   $S \in \text{InitConf } A \implies S \in \text{HAStates } A$   
apply (cut-tac A=A in InitConf-HAStates)  
apply fast  
done
```

```
lemma HAINitState-InitConf [simp]:  
  HAINitState A  $\in$  InitConf A  
by (unfold HAINitState-def InitConf-def, auto)
```

```
lemma InitConf-HAINitState-HARoot:  
   $\llbracket S \in \text{InitConf } A; S \neq \text{HAINitState } A \rrbracket \implies S \notin \text{States } (\text{HARoot } A)$   
apply (unfold InitConf-def)  
apply auto  
apply (rule mp)  
prefer 2  
apply fast  
back  
apply (rule mp)  
prefer 2  
apply fast  
back  
back  
apply (rule-tac b=S in rtrancl-induct)  
apply auto  
apply (simp add: ChiRel-HARoot)  
done
```

```
lemma InitConf-HARoot-HAINitState [simp]:  
   $\llbracket S \in \text{InitConf } A; S \in \text{States } (\text{HARoot } A) \rrbracket \implies S = \text{HAINitState } A$   
apply (subst not-not [THEN sym])  
apply (rule notI)  
apply (simp add:InitConf-HAINitState-HARoot)  
done
```

```

lemma HAINitState-CompFun-InitConf [simp]:
  [| SA ∈ the (CompFun A (HAINitState A)) |] ==> (InitState SA) ∈ InitConf A
apply (unfold InitConf-def HASTates-def)
apply auto
apply (rule rtrancl-Int)
apply auto
apply (cut-tac A=A and S=HAINitState A in HASTates-CompFun-States-ChiRel)
apply auto
apply (rule Image-singleton-iff [THEN subst])
apply (rotate-tac -1)
apply (drule sym)
apply simp
apply (rule-tac x=SA in beXI)
apply auto
done

```

```

lemma InitState-CompFun-InitConf:
  [| S ∈ HASTates A; SA ∈ the (CompFun A S); S ∈ InitConf A |] ==> (InitState
  SA) ∈ InitConf A
apply (unfold InitConf-def)
apply auto
apply (rule-tac b=S in rtrancl-into-rtrancl)
apply fast
apply (frule rtrancl-Int1)
apply auto
apply (case-tac S = HAINitState A)
apply simp
apply (rule rtrancl-mem-Sigma)
apply auto
apply (cut-tac A=A and S=S in HASTates-CompFun-States-ChiRel)
apply auto
apply (rule Image-singleton-iff [THEN subst])
apply (rotate-tac -1)
apply (drule sym)
apply simp
apply (rule-tac x=SA in beXI)
apply auto
done

```

```

lemma InitConf-HAINitStates:
  InitConf A ⊆ HAINitStates A
apply (unfold InitConf-def)
apply (rule subsetI)
apply auto
apply (frule rtrancl-Int1)
apply (case-tac x = HAINitState A)
apply simp
apply (rule rtrancl-mem-Sigma)
apply auto

```

done

lemma *InitState-notmem-InitConf*:

$\llbracket SA \in \text{the } (\text{CompFun } A \ S); S \in \text{InitConf } A; T \in \text{States } SA;$
 $T \neq \text{InitState } SA \rrbracket \implies T \notin \text{InitConf } A$

apply (*frule InitConf-HAStates2*)

apply (*unfold InitConf-def*)

apply *auto*

apply (*rule mp*)

prefer 2

apply *fast*

apply (*rule mp*)

prefer 2

apply *fast*

back

apply (*rule mp*)

prefer 2

apply *fast*

back

back

apply (*rule mp*)

prefer 2

apply *fast*

back

back

back

apply (*rule mp*)

prefer 2

apply *fast*

back

back

back

back

apply (*rule mp*)

prefer 2

apply *fast*

back

back

back

back

back

apply (*rule-tac b=T in rtrancl-induct*)

apply *auto*

done

lemma *InitConf-CompFun-InitState [simp]*:

$\llbracket SA \in \text{the } (\text{CompFun } A \ S); S \in \text{InitConf } A; T \in \text{States } SA;$
 $T \in \text{InitConf } A \rrbracket \implies T = \text{InitState } SA$

apply (*subst not-not [THEN sym]*)


```

apply (rule notI)
apply (frule InitState-notmem-InitConf)
apply auto
done

```

```

lemma InitConf-ChiRel-Ancestor:
   $\llbracket T \in \text{InitConf } A; (S, T) \in \text{ChiRel } A \rrbracket \implies S \in \text{InitConf } A$ 
apply (unfold InitConf-def)
apply auto
apply (erule rtranclE)
apply auto
apply (rename-tac U)
apply (cut-tac A=A in HAINitState-notmem-Range-ChiRel)
apply auto
apply (case-tac U = S)
apply (auto simp add: ChiRel-OneAncestor)
done

```

```

lemma InitConf-CompFun-Ancestor:
   $\llbracket S \in \text{HAStates } A; SA \in \text{the } (\text{CompFun } A \ S); T \in \text{InitConf } A; T \in \text{States } SA \rrbracket$ 
   $\implies S \in \text{InitConf } A$ 
apply (rule InitConf-ChiRel-Ancestor)
apply auto
apply (rule CompFun-ChiRel)
apply auto
done

```

7.2.16 StepConf

```

lemma StepConf-EmptySet [simp]:
   $\text{StepConf } A \ C \ \{\} = C$ 
by (unfold StepConf-def, auto)

end

```

8 Semantics of Hierarchical Automata

```

theory HASem
imports HA
begin

```

8.1 Definitions

```

definition
  RootExSem ::  $\llbracket ((\text{'s, 'e, 'd}) \text{seqauto}) \ \text{set}, \text{'s} \rightarrow (\text{'s, 'e, 'd}) \text{seqauto} \ \text{set},$ 
   $\text{'s} \ \text{set} \rrbracket \Rightarrow \text{bool}$  where
   $\text{RootExSem } F \ G \ C == (\exists! \ S. S \in \text{States } (\text{Root } F \ G) \wedge S \in C)$ 

```

```

definition

```

$UniqueSucStates :: [((\text{'s}, \text{'e}, \text{'d})seqauto) \text{ set}, \text{'s} \rightarrow (\text{'s}, \text{'e}, \text{'d})seqauto \text{ set}, \text{'s} \text{ set}] \Rightarrow \text{bool}$ **where**
 $UniqueSucStates \ F \ G \ C == \forall S \in (\bigcup (\text{States} \text{' } F)).$
 $\quad \forall A \in \text{the } (G \ S).$
 $\quad \text{if } (S \in C) \text{ then}$
 $\quad \quad \exists! S' . S' \in \text{States } A \wedge S' \in C$
 $\quad \text{else}$
 $\quad \quad \forall S \in \text{States } A. S \notin C$

definition

$IsConfSet :: [((\text{'s}, \text{'e}, \text{'d})seqauto) \text{ set}, \text{'s} \rightarrow (\text{'s}, \text{'e}, \text{'d})seqauto \text{ set}, \text{'s} \text{ set}] \Rightarrow \text{bool}$ **where**
 $IsConfSet \ F \ G \ C ==$
 $\quad C \subseteq (\bigcup (\text{States} \text{' } F)) \ \&$
 $\quad \text{RootExSem } F \ G \ C \ \&$
 $\quad UniqueSucStates \ F \ G \ C$

definition

$Status :: [(\text{'s}, \text{'e}, \text{'d})hierauto, \text{'s} \text{ set}, \text{'e} \text{ set}, \text{'d} \text{ data}] \Rightarrow \text{bool}$ **where**
 $Status \ HA \ C \ E \ D == E \subseteq \text{HAEvents } HA \wedge$
 $\quad IsConfSet \ (\text{SAs } HA) \ (\text{CompFun } HA) \ C \wedge$
 $\quad \text{Data.DataSpace } (\text{HAInitValue } HA) = \text{Data.DataSpace } D$

8.1.1 Status

lemma Status-EmptySet:

$(\text{Abs-hierauto } ((@ \ x . \ \text{True}),$
 $\quad \{\text{Abs-seqauto } (\{ @ \ x . \ \text{True}\}, (@ \ x . \ \text{True}), \{\}, \{\}), \{\}, \text{Map.empty}(@ \ x . \ \text{True}$
 $\mapsto \{\}),$
 $\quad \{ @ \ x . \ \text{True}\}, \{\}, @ \ x . \ \text{True}) \in$
 $\quad \{(HA, C, E, D) \mid HA \ C \ E \ D. \ \text{Status } HA \ C \ E \ D\}$
apply (*unfold Status-def CompFun-def SAs-def*)
apply *auto*
apply (*subst Abs-hierauto-inverse*)
apply (*subst hierauto-def*)
apply (*rule HierAuto-EmptySet*)
apply (*subst Abs-hierauto-inverse*)
apply (*subst hierauto-def*)
apply (*rule HierAuto-EmptySet*)
apply *auto*
apply (*unfold IsConfSet-def UniqueSucStates-def RootExSem-def*)
apply *auto*
apply (*unfold States-def*)
apply *auto*
apply (*unfold Root-def*)
apply (*rule someI2*)

```

apply (rule conjI)
apply fast
apply (simp add: ran-def)
apply simp
apply (subst Abs-seqauto-inverse)
apply (subst seqauto-def)
apply (rule SeqAuto-EmptySet)
apply simp
apply (unfold HAINitValue-def)
apply auto
apply (subst Abs-hierauto-inverse)
apply (subst hierauto-def)
apply (rule HierAuto-EmptySet)
apply simp
done

```

definition

```

status =
  {(HA, C, E, D) |
    (HA::('s, 'e, 'd) hierauto)
    (C::('s set))
    (E::('e set))
    (D::'d data). Status HA C E D}

```

```

typedef ('s, 'e, 'd) status =
  status :: (('s, 'e, 'd) hierauto * 's set * 'e set * 'd data) set
unfolding status-def
apply (rule exI)
apply (rule Status-EmptySet)
done

```

definition

```

HA :: ('s, 'e, 'd) status => ('s, 'e, 'd) hierauto where
HA == fst o Rep-status

```

definition

```

Conf :: ('s, 'e, 'd) status => 's set where
Conf == fst o snd o Rep-status

```

definition

```

Events :: ('s, 'e, 'd) status => 'e set where
Events == fst o snd o snd o Rep-status

```

definition

```

Value :: ('s, 'e, 'd) status => 'd data where
Value == snd o snd o snd o Rep-status

```

definition

$RootState :: ('s, 'e, 'd) status \Rightarrow 's$ **where**
 $RootState ST == @ S. S \in Conf ST \wedge S \in States (HARoot (HA ST))$

definition

$EnabledTrans :: (('s, 'e, 'd)status * ('s, 'e, 'd)seqauto * ('s, 'e, 'd)trans) set$ **where**
 $EnabledTrans == \{(ST, SA, T) .$
 $SA \in SAs (HA ST) \wedge$
 $T \in Delta SA \wedge$
 $source T \in Conf ST \wedge$
 $(Conf ST, Events ST, Value ST) \models (label T) \}$

definition

$ET :: ('s, 'e, 'd) status \Rightarrow (('s, 'e, 'd) trans) set$ **where**
 $ET ST == \bigcup SA \in SAs (HA ST). (EnabledTrans \text{ `` } \{ST\} \text{ `` } \{SA\})$

definition

$MaxNonConflict :: (('s, 'e, 'd)status,$
 $('s, 'e, 'd)trans set) \Rightarrow bool$ **where**
 $MaxNonConflict ST T ==$
 $(T \subseteq ET ST) \wedge$
 $(\forall A \in SAs (HA ST). card (T Int Delta A) \leq 1) \wedge$
 $(\forall t \in (ET ST). (t \in T) = (\neg (\exists t' \in ET ST. HigherPriority (HA ST)$
 $(t', t))))$

definition

$ResolveRacing :: ('s, 'e, 'd)trans set$
 $\Rightarrow ('d update set)$ **where**
 $ResolveRacing TS ==$
 let
 $U = PUpdate (Label TS)$
 in
 $SequentialRacing U$

definition

$HPT :: ('s, 'e, 'd)status \Rightarrow (('s, 'e, 'd)trans\ set)\ set$ **where**
 $HPT\ ST == \{ T. MaxNonConflict\ ST\ T\}$

definition

$InitStatus :: ('s, 'e, 'd)hierauto \Rightarrow ('s, 'e, 'd)status$ **where**
 $InitStatus\ A ==$
 $Abs-status\ (A, InitConf\ A, \{\}, HAINitValue\ A)$

definition

$StepActEvent :: ('s, 'e, 'd)trans\ set \Rightarrow 'e\ set$ **where**
 $StepActEvent\ TS == Union\ (Actevent\ (Label\ TS))$

definition

$StepStatus :: [('s, 'e, 'd)status, ('s, 'e, 'd)trans\ set, 'd\ update]$
 $\Rightarrow ('s, 'e, 'd)status$ **where**
 $StepStatus\ ST\ TS\ U =$
 $(let$
 $(A, C, E, D) = Rep-status\ ST;$
 $C' = StepConf\ A\ C\ TS;$
 $E' = StepActEvent\ TS;$
 $D' = U\ !!!\ D$
 in
 $Abs-status\ (A, C', E', D'))$

definition

$StepRelSem :: ('s, 'e, 'd)hierauto$
 $\Rightarrow (('s, 'e, 'd)status * ('s, 'e, 'd)status)\ set$ **where**
 $StepRelSem\ A == \{(ST, ST'). (HA\ ST) = A \wedge$
 $((HPT\ ST \neq \{\}) \rightarrow$

$$\begin{aligned}
& (\exists TS \in HPT\ ST. \\
& \quad \exists U \in ResolveRacing\ TS. \\
& \quad \quad ST' = StepStatus\ ST\ TS\ U) \ \& \\
& ((HPT\ ST = \{\}) \longrightarrow \\
& \quad (ST' = StepStatus\ ST\ \{\}\ DefaultUpdate))\}
\end{aligned}$$

inductive-set

ReachStati :: ('s,'e,'d)hierauto => ('s,'e,'d) status set
for *A* :: ('s,'e,'d)hierauto
where
Status0 : InitStatus *A* ∈ *ReachStati A*
| *StatusStep* :
 [[*ST* ∈ *ReachStati A*; *TS* ∈ *HPT ST*; *U* ∈ *ResolveRacing TS*]]
 ⇒ *StepStatus ST TS U* ∈ *ReachStati A*

8.2 Lemmas

lemma *Rep-status-tuple*:

Rep-status ST = (*HA ST*, *Conf ST*, *Events ST*, *Value ST*)
by (*unfold HA-def Conf-def Events-def Value-def, simp*)

lemma *Rep-status-select*:

(*HA ST*, *Conf ST*, *Events ST*, *Value ST*) ∈ *status*
by (*rule Rep-status-tuple [THEN subst], rule Rep-status*)

lemma *Status-select [simp]*:

Status (HA ST) (Conf ST) (Events ST) (Value ST)
apply (*cut-tac Rep-status-select*)
apply (*unfold status-def*)
apply *simp*
done

8.2.1 *IsConfSet*

lemma *IsConfSet-Status [simp]*:

IsConfSet (SAs (HA ST)) (CompFun (HA ST)) (Conf ST)
apply (*cut-tac Rep-status-select*)
apply (*unfold status-def Status-def*)
apply *auto*
done

8.2.2 *InitStatus*

```
lemma IsConfSet-InitConf [simp]:  
  IsConfSet (SAs A) (CompFun A) (InitConf A)  
apply (unfold IsConfSet-def RootExSem-def UniqueSucStates-def, fold HARoot-def)  
apply (rule conjI)  
apply (fold HASTates-def, simp)  
apply (rule conjI)  
apply (rule-tac a=HAINitState A in ex1I)  
apply auto  
apply (rename-tac S SA)  
apply (case-tac S ∈ InitConf A)  
apply auto  
apply (rule-tac x=InitState SA in exI)  
apply auto  
apply (rule InitState-CompFun-InitConf)  
apply auto  
apply (rename-tac S SA T U)  
apply (case-tac U = InitState SA)  
apply auto  
apply (simp only:InitConf-CompFun-Ancestor HASTates-SA-mem, simp)+  
done
```

```
lemma InitConf-status [simp]:  
  (A, InitConf A, {}, HAINitValue A) ∈ status  
apply (cut-tac Rep-status-select)  
apply (unfold status-def Status-def)  
apply auto  
done
```

```
lemma Conf-InitStatus-InitConf [simp]:  
  Conf (InitStatus A) = InitConf A  
apply (unfold Conf-def InitStatus-def)  
apply simp  
apply (subst Abs-status-inverse)  
apply auto  
done
```

```
lemma HAINitValue-Value-DataSpace-Status [simp]:  
  Data.DataSpace (HAINitValue (HA ST)) = Data.DataSpace (Value ST)  
apply (cut-tac Rep-status-select)  
apply (unfold status-def Status-def)  
apply fast  
done
```

```
lemma Value-InitStatus-HAINitValue [simp]:  
  Value (InitStatus A) = HAINitValue A  
apply (unfold Value-def InitStatus-def)  
apply simp  
apply (subst Abs-status-inverse)
```

apply *auto*
done

lemma *HA-InitStatus [simp]*:
 $HA (InitStatus A) = A$
apply (*unfold InitStatus-def HA-def*)
apply *auto*
apply (*subst Abs-status-inverse*)
apply *auto*
done

8.2.3 Events

lemma *Events-HAEvents-Status*:
 $(Events ST) \subseteq HAEvents (HA ST)$
apply (*cut-tac Rep-status-select*)
apply (*unfold status-def Status-def*)
apply *fast*
done

lemma *TS-EventSet*:
 $TS \subseteq ET ST \implies \bigcup (Actevent (Label TS)) \subseteq HAEvents (HA ST)$
apply (*unfold Actevent-def actevent-def ET-def EnabledTrans-def Action-def Label-def*)
apply (*cut-tac HA=HA ST in HAEvents-SAEvents-SAs*)
apply *auto*
apply (*rename-tac Event Source Trigger Guard Action Update Target*)
apply (*unfold SAEvents-def*)
apply (*erule subsetCE*)
apply *auto*
apply (*rename-tac SA*)
apply (*erule subsetCE*)
apply *auto*
apply (*erule-tac x=SA in ballE*)
apply *auto*
apply (*erule-tac x=(Trigger, Guard, Action, Update) in ballE*)
apply *auto*
apply (*cut-tac SA=SA in Label-Delta-subset*)
apply (*erule subsetCE*)
apply (*unfold Label-def image-def*)
apply *auto*
done

8.2.4 StepStatus

lemma *StepStatus-empty*:
 $Abs-status (HA ST, Conf ST, \{\}, U !!! (Value ST)) = StepStatus ST \{\} U$
apply (*unfold StepStatus-def Let-def*)
apply *auto*
apply (*subst Rep-status-tuple*)

apply *auto*
apply (*unfold StepActEvent-def*)
apply *auto*
done

lemma *status-empty-eventset [simp]*:
 $(HA\ ST, Conf\ ST, \{\}, U\ !!!\ (Value\ ST)) \in status$
apply (*unfold status-def Status-def*)
apply *auto*
done

lemma *HA-StepStatus-emptyTS [simp]*:
 $HA\ (StepStatus\ ST\ \{\}\ U) = HA\ ST$
apply (*subst StepStatus-empty [THEN sym]*)
apply (*unfold HA-def*)
apply *auto*
apply (*subst Abs-status-inverse*)
apply *auto*
apply (*subst Rep-status-tuple*)
apply *auto*
done

8.2.5 Enabled Transitions *ET*

lemma *HPT-ETI*:
 $TS \in HPT\ ST \implies TS \subseteq ET\ ST$
by (*unfold HPT-def MaxNonConflict-def, auto*)

lemma *finite-ET [simp]*:
 $finite\ (ET\ ST)$
by (*unfold ET-def Image-def EnabledTrans-def, auto*)

8.2.6 Finite Transition Set

lemma *finite-MaxNonConflict [simp]*:
 $MaxNonConflict\ ST\ TS \implies finite\ TS$
apply (*unfold MaxNonConflict-def*)
apply *auto*
apply (*subst finite-subset*)
apply *auto*
done

lemma *finite-HPT [simp]*:
 $TS \in HPT\ ST \implies finite\ TS$
by (*unfold HPT-def, auto*)

8.2.7 *PUpdate*

lemma *finite-Update*:
 $finite\ TS \implies finite\ ((\lambda F. (Rep-pupdate\ F)\ (Value\ ST))\ ' (PUpdate\ (Label\ TS)))$

by (rule finite-imageI, auto)

lemma *finite-PUupdate*:
 $TS \in \text{HPT } S \implies \text{finite } (\text{Expr.PUupdate } (\text{Label } TS))$
 apply auto
 done

lemma *HPT-ResolveRacing-Some* [simp]:
 $TS \in \text{HPT } S \implies (\text{SOME } u. u \in \text{ResolveRacing } TS) \in \text{ResolveRacing } TS$
 apply (unfold ResolveRacing-def Let-def)
 apply (rule finite-SequentialRacing)
 apply auto
 done

8.2.8 Higher Priority Transitions *HPT*

lemma *finite-HPT2* [simp]:
 $\text{finite } (\text{HPT } ST)$
 apply (cut-tac $ST=ST$ in finite-ET)
 apply (unfold HPT-def MaxNonConflict-def)
 apply (subst Collect-subset)
 apply (frule finite-Collect-subsets)
 apply auto
 done

lemma *HPT-target-StepConf* [simp]:
 $\llbracket TS \in \text{HPT } ST; T \in TS \rrbracket \implies \text{target } T \in \text{StepConf } (\text{HA } ST) (\text{Conf } ST) TS$
 apply (unfold StepConf-def)
 apply auto
 done

lemma *HPT-target-StepConf2* [simp]:
 $\llbracket TS \in \text{HPT } ST; (S,L,T) \in TS \rrbracket \implies T \in \text{StepConf } (\text{HA } ST) (\text{Conf } ST) TS$
 apply (unfold StepConf-def Target-def Source-def source-def target-def image-def)
 apply auto
 apply auto
 done

8.2.9 Delta Transition Set

lemma *ET-Delta*:
 $\llbracket TS \subseteq \text{ET } ST; t \in TS; \text{source } t \in \text{States } A; A \in \text{SAs } (\text{HA } ST) \rrbracket \implies t \in \text{Delta } A$
 apply (unfold ET-def EnabledTrans-def)
 apply simp
 apply (erule subsetCE)
 apply auto
 apply (rename-tac SA)
 apply (case-tac $A = SA$)
 apply auto

apply (*cut-tac HA=HA ST in MutuallyDistinct-HA*)
apply (*unfold MutuallyDistinct-def*)
apply *force*
done

lemma *ET-Delta-target:*

$\llbracket TS \subseteq ET\ ST; t \in TS; \text{target } t \in \text{States } A; A \in \text{SAs } (HA\ ST) \rrbracket \implies t \in \text{Delta}$
 A
apply (*unfold ET-def EnabledTrans-def*)
apply *simp*
apply (*erule subsetCE*)
apply *auto*
apply (*rename-tac SA*)
apply (*case-tac A = SA*)
apply *auto*
apply (*cut-tac HA=HA ST in MutuallyDistinct-HA*)
apply (*unfold MutuallyDistinct-def*)
apply *force*
done

lemma *ET-HADelta:*

$\llbracket TS \subseteq ET\ ST; t \in TS \rrbracket \implies t \in \text{HADelta } (HA\ ST)$
apply (*unfold HADelta-def*)
apply *auto*
apply (*unfold ET-def EnabledTrans-def Image-def*)
apply *auto*
done

lemma *HPT-HADelta:*

$\llbracket TS \in \text{HPT } ST; t \in TS \rrbracket \implies t \in \text{HADelta } (HA\ ST)$
apply (*rule ET-HADelta*)
apply (*unfold HPT-def MaxNonConflict-def*)
apply *auto*
done

lemma *HPT-Delta:*

$\llbracket TS \in \text{HPT } ST; t \in TS; \text{source } t \in \text{States } A; A \in \text{SAs } (HA\ ST) \rrbracket \implies t \in \text{Delta}$
 A
apply (*rule ET-Delta*)
apply *auto*
apply (*unfold HPT-def MaxNonConflict-def*)
apply *fast*
done

lemma *HPT-Delta-target:*

$\llbracket TS \in \text{HPT } ST; t \in TS; \text{target } t \in \text{States } A; A \in \text{SAs } (HA\ ST) \rrbracket \implies t \in \text{Delta}$
 A
apply (*rule ET-Delta-target*)
apply *auto*

apply (*unfold HPT-def MaxNonConflict-def*)
apply *fast*
done

lemma *OneTrans-HPT-SA*:
 $\llbracket TS \in HPT\ ST; T \in TS; \text{source } T \in \text{States } SA; U \in TS; \text{source } U \in \text{States } SA; SA \in SAs\ (HA\ ST) \rrbracket \implies T = U$
apply (*unfold HPT-def MaxNonConflict-def Source-def*)
apply *auto*
apply (*erule-tac x=SA in ballE*)
apply (*case-tac finite (TS \cap Delta SA)*)
apply (*frule-tac t=T in OneElement-Card*)
apply *fast*
apply (*frule-tac t=T and A=SA in ET-Delta*)
apply *assumption+*
apply *fast*
apply (*frule-tac t=U in OneElement-Card*)
apply *fast*
apply (*frule-tac t=U and A=SA in ET-Delta*)
apply *auto*
done

lemma *OneTrans-HPT-SA2*:
 $\llbracket TS \in HPT\ ST; T \in TS; \text{target } T \in \text{States } SA; U \in TS; \text{target } U \in \text{States } SA; SA \in SAs\ (HA\ ST) \rrbracket \implies T = U$
apply (*unfold HPT-def MaxNonConflict-def Target-def*)
apply *auto*
apply (*erule-tac x=SA in ballE*)
apply (*case-tac finite (TS \cap Delta SA)*)
apply (*frule-tac t=T in OneElement-Card*)
apply *fast*
apply (*frule-tac t=T and A=SA in ET-Delta-target*)
apply *assumption+*
apply *fast*
apply (*frule-tac t=U in OneElement-Card*)
apply *fast*
apply (*frule-tac t=U and A=SA in ET-Delta-target*)
apply *auto*
done

8.2.10 Target Transition Set

lemma *ET-Target-HAStates*:
 $TS \subseteq ET\ ST \implies \text{Target } TS \subseteq HAStates\ (HA\ ST)$
apply (*unfold HAStates-def Target-def target-def ET-def EnabledTrans-def Action-def Label-def*)
apply (*cut-tac HA=HA ST in Target-SAs-Delta-States*)
apply *auto*
apply (*rename-tac Source Trigger Guard Action Update Target*)

```

apply (unfold Target-def)
apply (erule subsetCE)
apply auto
apply (rename-tac SA)
apply (erule subsetCE)
apply auto
apply (unfold image-def)
apply auto
apply (metis target-select)
done

```

```

lemma HPT-Target-HAStates:
   $TS \in \text{HPT } ST \implies \text{Target } TS \subseteq \text{HAStates } (HA \ ST)$ 
apply (rule HPT-ETI [THEN ET-Target-HAStates])
apply assumption
done

```

```

lemma HPT-Target-HAStates2 [simp]:
   $\llbracket TS \in \text{HPT } ST; S \in \text{Target } TS \rrbracket \implies S \in \text{HAStates } (HA \ ST)$ 
apply (cut-tac HPT-Target-HAStates)
apply fast+
done

```

```

lemma OneState-HPT-Target:
   $\llbracket TS \in \text{HPT } ST; S \in \text{Target } TS;$ 
   $T \in \text{Target } TS; S \in \text{States } SA;$ 
   $T \in \text{States } SA; SA \in \text{SAs } (HA \ ST) \rrbracket$ 
   $\implies S = T$ 
apply (unfold Target-def)
apply (auto dest: OneTrans-HPT-SA2[rotated -1])
done

```

8.2.11 Source Transition Set

```

lemma ET-Source-Conf:
   $TS \subseteq \text{ET } ST \implies (\text{Source } TS) \subseteq \text{Conf } ST$ 
apply (unfold Source-def ET-def EnabledTrans-def)
apply auto
done

```

```

lemma HPT-Source-Conf [simp]:
   $TS \in \text{HPT } ST \implies (\text{Source } TS) \subseteq \text{Conf } ST$ 
apply (unfold HPT-def MaxNonConflict-def)
apply (rule ET-Source-Conf)
apply auto
done

```

```

lemma ET-Source-Target [simp]:
   $\llbracket SA \in \text{SAs } (HA \ ST); TS \subseteq \text{ET } ST; \text{States } SA \cap \text{Source } TS = \{\} \rrbracket \implies \text{States}$ 

```

```

SA  $\cap$  Target TS = {}
apply (unfold ET-def EnabledTrans-def Source-def Target-def)
apply auto
apply (rename-tac Source Trigger Guard Action Update Target)
apply (erule subsetCE)
apply auto
apply (rename-tac SAA)
apply (unfold image-def source-def Int-def)
apply auto
apply (erule-tac x=Source in allE)
apply auto
apply (frule Delta-source-States)
apply (unfold source-def)
apply auto
apply (case-tac SA=SAA)
apply auto
apply (cut-tac HA=HA ST in MutuallyDistinct-HA)
apply (unfold MutuallyDistinct-def)
apply (erule-tac x=SA in ballE)
apply (erule-tac x=SAA in ballE)
apply auto
apply (frule Delta-target-States)
apply (unfold target-def)
apply force
done

```

lemma *HPT-Source-Target* [simp]:
 $\llbracket TS \in \text{HPT } ST; \text{States } SA \cap \text{Source } TS = \{\}; SA \in \text{SAs } (HA \text{ } ST) \rrbracket \implies \text{States } SA \cap \text{Target } TS = \{\}$
apply (unfold HPT-def MaxNonConflict-def)
apply auto
done

lemma *ET-target-source*:
 $\llbracket TS \subseteq \text{ET } ST; t \in TS; \text{target } t \in \text{States } A; A \in \text{SAs } (HA \text{ } ST) \rrbracket \implies \text{source } t \in \text{States } A$
apply (frule ET-Delta-target)
apply auto
done

lemma *ET-source-target*:
 $\llbracket TS \subseteq \text{ET } ST; t \in TS; \text{source } t \in \text{States } A; A \in \text{SAs } (HA \text{ } ST) \rrbracket \implies \text{target } t \in \text{States } A$
apply (frule ET-Delta)
apply auto
done

lemma *HPT-target-source*:
 $\llbracket TS \in \text{HPT } ST; t \in TS; \text{target } t \in \text{States } A; A \in \text{SAs } (HA \text{ } ST) \rrbracket \implies \text{source } t$

$\in \text{States } A$
apply (rule *ET-target-source*)
apply *auto*
apply (unfold *HPT-def MaxNonConflict-def*)
apply *fast*
done

lemma *HPT-source-target*:
 $\llbracket TS \in \text{HPT } ST; t \in \text{States } A; A \in \text{SAs } (HA \text{ } ST) \rrbracket \implies \text{target } t$
 $\in \text{States } A$
apply (rule *ET-source-target*)
apply *auto*
apply (unfold *HPT-def MaxNonConflict-def*)
apply *fast*
done

lemma *HPT-source-target2 [simp]*:
 $\llbracket TS \in \text{HPT } ST; (s,l,t) \in TS; s \in \text{States } A; A \in \text{SAs } (HA \text{ } ST) \rrbracket \implies t \in \text{States } A$
apply (cut-tac *ST=ST and TS=TS and t=(s,l,t) in HPT-source-target*)
apply *auto*
done

lemma *ChiRel-ChiStar-Source-notmem*:
 $\llbracket TS \in \text{HPT } ST; (S, T) \in \text{ChiRel } (HA \text{ } ST); S \in \text{Conf } ST; T \notin \text{ChiStar } (HA \text{ } ST) \text{ ``Source } TS \rrbracket \implies S \notin \text{ChiStar } (HA \text{ } ST) \text{ ``Source } TS$
apply *auto*
apply (rename-tac *U*)
apply (simp only: *Image-def*)
apply *auto*
apply (erule-tac *x=U in ballE*)
apply (fast intro: *ChiRel-ChiStar-trans*)+
done

lemma *ChiRel-ChiStar-notmem*:
 $\llbracket TS \in \text{HPT } ST; (S, T) \in \text{ChiRel } (HA \text{ } ST); S \in \text{ChiStar } (HA \text{ } ST) \text{ ``Source } TS \rrbracket \implies T \notin \text{Source } TS$
using $\llbracket \text{hypsubst-thin} = \text{true} \rrbracket$
apply (unfold *HPT-def MaxNonConflict-def HigherPriority-def restrict-def*)
apply *auto*
apply (rename-tac *U*)
apply (unfold *Source-def image-def*)
apply *auto*
apply (rename-tac *SSource STrigger SGuard SAction SUpdate STarget TSource TTrigger TGuard TAction TUpdate TTarget*)
apply (erule-tac *x=(SSource, (STrigger, SGuard, SAction, SUpdate), STarget) in ballE*)
apply *auto*

```

apply (erule-tac x=(TSource, (TTrigger, TGuard, TAction, TUpdate), TTarget)
in ballE)
apply auto
apply (simp add: ET-HADelta)
apply (case-tac SSource=S)
apply auto
apply (frule ChiStar-ChiPlus-noteq)
apply fast
apply (fast intro: ChiRel-ChiPlus-trans)
done

```

8.2.12 StepActEvents

```

lemma StepActEvent-empty [simp]:
  StepActEvent {} = {}
by (unfold StepActEvent-def, auto)

```

```

lemma StepActEvent-HAEvents:
  TS ∈ HPT ST ⇒ StepActEvent TS ⊆ HAEvents (HA ST)
apply (unfold StepActEvent-def image-def)
apply (rule HPT-ETI [THEN TS-EventSet])
apply assumption
done

```

8.2.13 UniqueSucStates

```

lemma UniqueSucStates-Status [simp]:
  UniqueSucStates (SAs (HA ST)) (CompFun (HA ST)) (Conf ST)
apply (cut-tac Rep-status-select)
apply (unfold status-def Status-def IsConfSet-def)
apply auto
done

```

8.2.14 RootState

```

lemma RootExSem-Status [simp]:
  RootExSem (SAs (HA ST)) (CompFun (HA ST)) (Conf ST)
apply (cut-tac Rep-status-select)
apply (unfold status-def Status-def IsConfSet-def)
apply auto
done

```

```

lemma RootState-HARootState [simp]:
  (RootState ST) ∈ States (HARoot (HA ST))
apply (unfold RootState-def)
apply (cut-tac ST=ST in RootExSem-Status)
apply (unfold RootExSem-def HARoot-def HAStates-def)
apply auto
apply (subst some1-equality)
apply auto

```


done

lemma *RootState-Conf* [simp]:
 $(\text{RootState } ST) \in (\text{Conf } ST)$
apply (unfold *RootState-def*)
apply (cut-tac $ST=ST$ in *RootExSem-Status*)
apply (unfold *RootExSem-def* *HARoot-def* *HAStates-def*)
apply auto
apply (subst *some1-equality*)
apply auto
done

lemma *RootState-notmem-Chi* [simp]:
 $S \in \text{HAStates } (HA \ ST) \implies (\text{RootState } ST) \notin \text{Chi } (HA \ ST) \ S$
by auto

lemma *RootState-notmem-Range-ChiRel* [simp]:
 $\text{RootState } ST \notin \text{Range } (\text{ChiRel } (HA \ ST))$
by auto

lemma *RootState-Range-ChiPlus* [simp]:
 $\text{RootState } ST \notin \text{Range } (\text{ChiPlus } (HA \ ST))$
by auto

lemma *RootState-Range-ChiStar* [simp]:
 $\llbracket x \neq \text{RootState } ST \rrbracket \implies (x, \text{RootState } ST) \notin (\text{ChiStar } (HA \ ST))$
by auto

lemma *RootState-notmem-ChiRel* [simp]:
 $(x, \text{RootState } ST) \notin (\text{ChiRel } (HA \ ST))$
by (unfold *ChiRel-def*, auto)

lemma *RootState-notmem-ChiRel2* [simp]:
 $\llbracket S \in \text{States } (\text{HARoot } (HA \ ST)) \rrbracket \implies (x, S) \notin (\text{ChiRel } (HA \ ST))$
by (unfold *ChiRel-def*, auto)

lemma *RootState-Conf-StepConf* [simp]:
 $\llbracket \text{RootState } ST \notin \text{Source } TS \rrbracket \implies \text{RootState } ST \in \text{StepConf } (HA \ ST) (\text{Conf } ST) \ TS$
apply (unfold *StepConf-def*)
apply auto
apply (rename-tac *S*)
apply (case-tac $S=\text{RootState } ST$)
apply fast
apply auto
apply (rename-tac *S*)
apply (case-tac $S=\text{RootState } ST$)
apply fast
apply auto

done

lemma *OneRootState-Conf* [simp]:

$\llbracket S \in \text{States } (\text{HARoot } (\text{HA } ST)); S \in \text{Conf } ST \rrbracket \implies S = \text{RootState } ST$
apply (*cut-tac* $ST=ST$ **in** *IsConfSet-Status*)
apply (*unfold* *IsConfSet-def* *RootExSem-def*)
apply (*fold* *HARoot-def*)
apply *auto*
done

lemma *OneRootState-Source*:

$\llbracket TS \in \text{HPT } ST; S \in \text{Source } TS; S \in \text{States } (\text{HARoot } (\text{HA } ST)) \rrbracket \implies S = \text{RootState } ST$
apply (*cut-tac* $ST=ST$ **and** $TS=TS$ **in** *HPT-Source-Conf*, *assumption*)
apply (*cut-tac* $ST=ST$ **in** *OneRootState-Conf*)
apply *fast+*
done

lemma *OneState-HPT-Target-Source*:

$\llbracket TS \in \text{HPT } ST; S \in \text{States } SA; SA \in \text{SAs } (\text{HA } ST);$
 $\text{States } SA \cap \text{Source } TS = \{\} \rrbracket$
 $\implies S \notin \text{Target } TS$
apply (*unfold* *Target-def*)
apply *auto*
apply (*unfold* *Source-def* *Image-def* *Int-def*)
apply *auto*
apply (*frule* *HPT-target-source*)
apply *auto*
done

lemma *RootState-notmem-Target* [simp]:

$\llbracket TS \in \text{HPT } ST; S \in \text{States } (\text{HARoot } (\text{HA } ST)); \text{RootState } ST \notin \text{Source } TS \rrbracket$
 $\implies S \notin \text{Target } TS$
apply *auto*
apply (*frule* *OneState-HPT-Target-Source*)
prefer 4
apply *fast+*
apply *simp*
apply (*unfold* *Int-def*)
apply *auto*
apply (*frule* *OneRootState-Source*)
apply *fast+*
done

8.2.15 Configuration *Conf*

lemma *Conf-HAStates*:

$\text{Conf } ST \subseteq \text{HAStates } (\text{HA } ST)$
apply (*cut-tac* *Rep-status-select*)

apply (*unfold IsConfSet-def status-def Status-def HAStates-def*)
apply *fast*
done

lemma *Conf-HAStates2 [simp]*:
 $S \in \text{Conf } ST \implies S \in \text{HAStates } (HA \ ST)$
apply (*cut-tac ST=ST in Conf-HAStates*)
apply *fast*
done

lemma *OneState-Conf [intro]*:
 $\llbracket S \in \text{Conf } ST; T \in \text{Conf } ST; S \in \text{States } SA; T \in \text{States } SA; SA \in \text{SAs } (HA \ ST) \rrbracket \implies T = S$
apply (*cut-tac ST=ST in IsConfSet-Status*)
apply (*unfold IsConfSet-def UniqueSucStates-def*)
apply (*case-tac SA = HARoot (HA ST)*)
apply (*cut-tac ST=ST and S=S in OneRootState-Conf*)
apply *fast+*
apply (*simp only:OneRootState-Conf*)
apply (*erule conjE*)
apply (*cut-tac HA=HA ST in OneAncestor-HA*)
apply (*unfold OneAncestor-def*)
apply (*fold HARoot-def*)
apply (*erule-tac x=SA in ballE*)
apply (*erule ex1-implies-ex*)
apply (*erule exE*)
apply (*rename-tac U*)
apply (*erule-tac x=U in ballE*)
apply (*erule-tac x=SA in ballE*)
apply (*case-tac U ∈ Conf ST*)
apply *simp*
apply *safe*
apply *fast+*
apply *simp*
apply *fast*
done

lemma *OneState-HPT-SA*:
 $\llbracket TS \in \text{HPT } ST; S \in \text{Source } TS; T \in \text{Source } TS; S \in \text{States } SA; T \in \text{States } SA; SA \in \text{SAs } (HA \ ST) \rrbracket \implies S = T$
apply (*rule OneState-Conf*)
apply *auto*
apply (*frule HPT-Source-Conf, fast*)
done

lemma *HPT-SAStates-Target-Source*:
 $\llbracket TS \in \text{HPT } ST; A \in \text{SAs } (HA \ ST); S \in \text{States } A; T \in \text{States } A; S \in \text{Conf } ST; T \in \text{Target } TS \rrbracket \implies S \in \text{Source } TS$

```

apply (case-tac  $States\ A \cap Source\ TS = \{\}$ )
apply (frule OneState-HPT-Target-Source)
apply fast
back
apply simp+
apply auto
apply (rename-tac  $U$ )
apply (cut-tac  $ST=ST$  in HPT-Source-Conf)
apply fast
apply (frule-tac  $S=S$  and  $T=U$  in OneState-Conf)
apply fast+
done

```

```

lemma HPT-Conf-Target-Source:
   $\llbracket TS \in HPT\ ST; S \in Conf\ ST; S \in Target\ TS \rrbracket \implies S \in Source\ TS$ 
apply (frule Conf-HAStates2)
apply (unfold HAStates-def)
apply auto
apply (simp only:HPT-SAStates-Target-Source)
done

```

```

lemma Conf-SA:
   $S \in Conf\ ST \implies \exists A \in SAs\ (HA\ ST). S \in States\ A$ 
apply (cut-tac  $ST=ST$  in IsConfSet-Status)
apply (unfold IsConfSet-def)
apply fast
done

```

```

lemma HPT-Source-HAStates [simp]:
   $\llbracket TS \in HPT\ ST; S \in Source\ TS \rrbracket \implies S \in HAStates\ (HA\ ST)$ 
apply (frule HPT-Source-Conf)
apply (rule Conf-HAStates2)
apply fast
done

```

```

lemma Conf-Ancestor:
   $\llbracket S \in Conf\ ST; A \in the\ (CompFun\ (HA\ ST)\ S) \rrbracket \implies \exists! T \in States\ A. T \in Conf\ ST$ 
apply (cut-tac  $ST=ST$  in IsConfSet-Status)
apply (unfold IsConfSet-def UniqueSucStates-def)
apply safe
apply (erule-tac  $x=S$  in ballE)
prefer 2
apply blast
apply (erule-tac  $x=A$  in ballE)
prefer 2
apply fast
apply simp

```

apply (*fast intro: HAStates-CompFun-SAs-mem Conf-HAStates2*)
done

lemma *Conf-ChiRel:*

$\llbracket (S,T) \in \text{ChiRel } (HA \ ST); T \in \text{Conf } ST \rrbracket \implies S \in \text{Conf } ST$
apply (*unfold ChiRel-def Chi-def restrict-def*)
apply *simp*
apply *safe*
apply *simp*
apply *safe*
apply (*rename-tac SA*)
apply (*unfold HAStates-def*)
apply *simp*
apply *safe*
apply (*rename-tac U*)
apply (*cut-tac ST=ST in UniqueSucStates-Status*)
apply (*unfold UniqueSucStates-def*)
apply (*erule-tac x=S in ballE*)
apply (*erule-tac x=SA in ballE*)
apply *auto*
apply (*case-tac S ∈ Conf ST*)
apply *simp+*
done

lemma *Conf-ChiPlus:*

$\llbracket (T,S) \in \text{ChiPlus } (HA \ ST) \rrbracket \implies S \in \text{Conf } ST \longrightarrow T \in \text{Conf } ST$
apply (*unfold ChiPlus-def*)
apply (*rule-tac a=T and b=S and r=(ChiRel (HA ST)) in trancl-induct*)
apply (*fast intro: Conf-ChiRel*)
done

lemma *HPT-Conf-Target-Source-ChiPlus:*

$\llbracket TS \in \text{HPT } ST; S \in \text{Conf } ST; S \in \text{ChiPlus } (HA \ ST) \text{ “ Target } TS \rrbracket$
 $\implies S \in \text{ChiStar } (HA \ ST) \text{ “ Source } TS$
apply *auto*
apply (*rename-tac T*)
apply (*simp add: Image-def*)
apply (*frule HPT-Target-HAStates2*)
apply *fast*
apply (*unfold HAStates-def*)
apply *auto*
apply (*rename-tac SA*)
apply (*case-tac States SA ∩ Source TS = {}*)
apply (*simp only: OneState-HPT-Target-Source*)
apply *auto*
apply (*rename-tac U*)
apply (*erule-tac x=U in ballE*)
apply *auto*

```

apply (case-tac  $U=T$ )
apply auto
apply (frule Conf-ChiPlus)
apply simp
apply (frule HPT-Conf-Target-Source)
apply fast
back
apply fast
apply (simp add:OneState-HPT-SA)
done

lemma OneState-HPT-Target-ChiRel:
   $\llbracket TS \in \text{HPT } ST; (U, T) \in \text{ChiRel } (HA \text{ } ST);$ 
   $U \in \text{Target } TS; A \in \text{SAs } (HA \text{ } ST); T \in \text{States } A;$ 
   $S \in \text{States } A \rrbracket \implies S \notin \text{Target } TS$ 
using  $\llbracket \text{hypsubst-thin} = \text{true} \rrbracket$ 
apply auto
apply (unfold HigherPriority-def restrict-def HPT-def MaxNonConflict-def Target-def)
apply auto
apply (rename-tac SSource STrigger SGuard SAction SUpdate STarget
  TSource TTrigger TGuard TAction TUpdate TTarget)
apply (cut-tac  $t=(TSource, (TTrigger, TGuard, TAction, TUpdate), TTarget)$  and
   $TS=TS$  and  $ST=ST$  and  $A=A$  in ET-target-source)
apply assumption+
apply simp
apply assumption
apply (frule ChiRel-HAStates)
apply (unfold HAStates-def)
apply safe
apply (cut-tac  $t=(SSource, (STrigger, SGuard, SAction, SUpdate), STarget)$  and
   $A=x$  and  $ST=ST$  and  $TS=TS$  in ET-target-source)
apply assumption+
apply simp
apply assumption
apply simp
apply (erule-tac  $x=(SSource, (STrigger, SGuard, SAction, SUpdate), STarget)$  in
  ballE)
apply simp
apply (erule-tac  $x=(TSource, (TTrigger, TGuard, TAction, TUpdate), TTarget)$ 
  in ballE)
apply (simp add:ET-HADelta)
apply (cut-tac  $A=HA \text{ } ST$  and  $S=STarget$  and  $T=T$  and  $U=TSource$  in ChiRel-SA-OneAncestor)
apply fast+
apply (frule ET-Source-Conf)
apply (unfold Source-def image-def)
apply (case-tac  $SSource \in \text{Conf } ST$ )
prefer 2
apply (erule subsetCE)

```

```

back
apply fast
back
apply simp
apply (case-tac TSource ∈ Conf ST)
prefer 2
apply (erule subsetCE)
back
apply fast
apply simp
apply (case-tac STarget=SSource)
apply simp
apply (fast intro: Conf-ChiRel)+
done

```

```

lemma OneState-HPT-Target-ChiPlus [rule-format]:
  [[ TS ∈ HPT ST; (U,T) ∈ ChiPlus (HA ST);
    S ∈ Target TS; A ∈ SAs (HA ST);
    S ∈ States A ]] ⇒ T ∈ States A ⟶ U ∉ Target TS
using [[hypsubst-thin = true]]
apply (unfold ChiPlus-def)
apply (rule-tac a=U and b=T and r=(ChiRel (HA ST)) in converse-trancl-induct)
apply auto
apply (simp only: OneState-HPT-Target-ChiRel)
apply (rename-tac V W)
apply (fold ChiPlus-def)
apply (unfold HPT-def MaxNonConflict-def Target-def HigherPriority-def restrict-def)
apply auto
apply (rename-tac SSource STrigger SGuard SAction SUpdate STarget
  TSource TTrigger TGuard TAction TUpdate TTarget)
apply (cut-tac t=(SSource, (STrigger, SGuard, SAction, SUpdate), STarget) and
  ST=ST and TS=TS and A=A in ET-target-source)
apply assumption+
apply simp
apply assumption
apply simp
apply (frule ChiRel-HAStates)
apply (unfold HAStates-def)
apply safe
apply (cut-tac t=(TSource, (TTrigger, TGuard, TAction, TUpdate), TTarget) and
  A=x and TS=TS and ST=ST in ET-target-source)
apply assumption+
apply simp
apply assumption
apply simp
apply (erule-tac x=(TSource, (TTrigger, TGuard, TAction, TUpdate), TTarget)
  in ballE)
apply simp
apply (erule-tac x=(SSource, (STrigger, SGuard, SAction, SUpdate), STarget) in

```

```

ballE)
apply (simp add: ET-HADelta)
apply (cut-tac A=HA ST and S=TTTarget and T=T and U=SSource in ChiPlus-SA-OneAncestor)
apply (fast intro: ChiRel-ChiPlus-trans2)
apply fast+
apply (frule ET-Source-Conf)
apply (unfold Source-def image-def)
apply (case-tac SSource ∈ Conf ST)
prefer 2
apply (erule subsetCE)
back
apply fast
apply simp
apply (case-tac TSource ∈ Conf ST)
prefer 2
apply (erule subsetCE)
back
apply fast
back
apply simp
apply (case-tac TTTarget=SSource)
apply simp
apply (frule-tac T=TTarget and S=SSource in Conf-ChiPlus)
apply simp
apply (frule-tac T=TSource and S=TTarget in OneState-Conf)
apply fast+
done

```

8.2.16 RootExSem

lemma *RootExSem-StepConf*:

```

[[ TS ∈ HPT ST ]] ⇒
  RootExSem (SAs (HA ST)) (CompFun (HA ST)) (StepConf (HA ST) (Conf
ST) TS)
apply (unfold RootExSem-def)
apply (fold HARoot-def)
apply auto
apply (case-tac RootState ST ∉ Source TS)
apply (rule-tac x=RootState ST in exI)
apply simp
apply simp
apply (unfold Source-def image-def)
apply simp
apply (erule bexE)
apply (rename-tac T)
apply (rule-tac x=target T in exI)
apply simp
apply (rule HPT-source-target)
apply auto

```



```

apply (rename-tac  $S T$ )
apply (case-tac  $S \in \text{Conf } ST$ )
apply (case-tac  $T \in \text{Conf } ST$ )
apply (frule OneRootState-Conf)
apply auto
apply (frule OneRootState-Conf)
apply auto
apply (frule OneRootState-Conf)
apply auto
apply (case-tac  $\text{RootState } ST \in \text{Source } TS$ )
apply (case-tac  $T \in \text{Source } TS$ )
apply (frule HPT-Source-Conf)
apply fast
apply (unfold StepConf-def)
apply auto
apply (frule OneState-HPT-Target)
apply (frule-tac  $SA = \text{HARoot } (HA ST)$  and  $TS = TS$  and  $S = T$  and  $T = \text{RootState } ST$  in OneState-HPT-Target)
apply fast+
apply simp+
apply (frule trancl-Int-mem, fold ChiPlus-def, force)+
prefer 2
apply (frule OneState-HPT-Target)
apply fast+
back
apply simp+
apply (case-tac  $\text{RootState } ST \in \text{Source } TS$ )
apply (case-tac  $T = \text{RootState } ST$ )
apply auto
apply (frule trancl-Int-mem, fold ChiPlus-def, force)+
done

```

8.2.17 *StepConf*

lemma *Target-StepConf*:

```

 $S \in \text{Target } TS \implies S \in \text{StepConf } (HA ST) (\text{Conf } ST) TS$ 
apply (unfold StepConf-def)
apply auto
done

```

lemma *Target-ChiRel-HAInit-StepConf*:

```

 $\llbracket S \in \text{Target } TS; (S, T) \in \text{ChiRel } A; T \in \text{HAInitStates } A \rrbracket \implies T \in \text{StepConf } A C TS$ 
apply (unfold StepConf-def)
apply auto
done

```

lemma *StepConf-HAStates*:

```

 $TS \in \text{HPT } ST \implies \text{StepConf } (HA ST) (\text{Conf } ST) TS \subseteq \text{HAStates } (HA ST)$ 

```

apply (*unfold StepConf-def*)
apply *auto*
apply (*frule tranclD2*)
apply *auto*
done

lemma *RootState-Conf-StepConf2 [simp]*:
 $\llbracket \text{source } T = \text{RootState } ST; T \in TS \rrbracket \implies \text{target } T \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$
apply (*unfold StepConf-def*)
apply *auto*
done

lemma *HPT-StepConf-HAStates [simp]*:
 $\llbracket TS \in HPT \ ST; S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS \rrbracket \implies S \in HAStates \ (HA \ ST)$
apply (*unfold StepConf-def*)
apply *auto*
apply (*frule tranclD2*)
apply *auto*
done

lemma *StepConf-Target-HAInitStates*:
 $\llbracket S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS; S \notin \text{Target } TS; S \notin \text{Conf } ST \rrbracket \implies S \in HAInitStates \ (HA \ ST)$
apply (*unfold StepConf-def*)
apply *auto*
apply (*frule tranclD2*)
apply *auto*
done

lemma *InitSucState-StepConf*:
 $\llbracket TS \in HPT \ ST; S \notin \text{Target } TS; A \in \text{the } (CompFun \ (HA \ ST) \ S); S \notin \text{Conf } ST; S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS \rrbracket \implies$
 $\text{InitState } A \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$
apply (*frule StepConf-HAStates [THEN subsetD, THEN CompFun-HAInitStates-HAStates]*)
apply *fast+*
apply (*subst (asm) StepConf-def*)
apply *safe*
apply (*unfold StepConf-def*)
apply (*fast intro: HAInitStates-InitState-trancl*)
apply (*frule trancl-Int-mem, fold ChiPlus-def*)
apply (*fast intro: ChiPlus-HAStates-Right [THEN HAInitStates-InitState-trancl2]*)
done

lemma *InitSucState-Target-StepConf*:
 $\llbracket TS \in HPT \ ST; S \in \text{Target } TS; A \in \text{the } (CompFun \ (HA \ ST) \ S) \rrbracket \implies$
 $\text{InitState } A \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$

```

apply (frule HPT-Target-HAStates2 [THEN CompFun-HAInitStates-HAStates])
apply fast+
apply (frule HPT-Target-HAStates2 [THEN CompFun-ChiRel])
apply (fast intro:InitState-States)+
apply (unfold StepConf-def)
apply auto
done

```

lemma *InitSucState-Conf-StepConf*:

```

  [| TS ∈ HPT ST; S ∈ StepConf (HA ST) (Conf ST) TS;
    S ∉ Target TS; A ∈ the (CompFun (HA ST) S);
    S ∈ Conf ST; S ∈ ChiStar (HA ST) “(Source TS)” |] ⇒
    initState A ∈ StepConf (HA ST) (Conf ST) TS
apply (frule Conf-HAStates2 [THEN CompFun-HAInitStates-HAStates])
apply fast
apply (subst (asm) StepConf-def)
apply safe
apply fast
apply (unfold StepConf-def)
apply (fast intro:HAInitStates-InitState-trancl)
apply (rename-tac T U V)
apply (frule trancl-Int-mem, fold ChiPlus-def, safe)
apply (subst (asm) Image-def, safe)
apply (rule-tac x=U in be1)
apply (simp only: ChiPlus-HAStates-Right [THEN HAInitStates-InitState-trancl2])
apply fast
apply (subst (asm) Image-def, safe)
apply (rule-tac x=U in be1)
apply (simp only: ChiPlus-HAStates-Right [THEN HAInitStates-InitState-trancl2])
apply fast
done

```

lemma *SucState-Conf-StepConf*:

```

  [| TS ∈ HPT ST; S ∈ StepConf (HA ST) (Conf ST) TS;
    S ∉ Target TS; A ∈ the (CompFun (HA ST) S);
    S ∈ Conf ST; States A ∩ ChiStar (HA ST) “(Source TS)” = {} |] ⇒
    ∃ x. x ∈ States A ∧ x ∈ StepConf (HA ST) (Conf ST) TS
apply (unfold StepConf-def)
apply (cut-tac ST=ST in UniqueSucStates-Status)
apply (unfold UniqueSucStates-def)
apply (cut-tac ST=ST in Conf-HAStates2)
apply fast
apply (fold HAStates-def)
apply (erule-tac x=S in ballE)
apply (erule-tac x=A in ballE)
apply simp
apply fast+
done

```

lemma *SucState-Conf-Source-StepConf*:

```

[[ TS ∈ HPT ST; S ∈ StepConf (HA ST) (Conf ST) TS;
   S ∉ Target TS; A ∈ the (CompFun (HA ST) S);
   S ∈ Conf ST; States A ∩ ChiStar (HA ST) “ (Source TS) ≠ {};
   S ∉ ChiStar (HA ST) “ (Source TS) ]] ⇒
  ∃ x. x ∈ States A ∧ x ∈ StepConf (HA ST) (Conf ST) TS
apply safe
apply (rename-tac T U)
apply (frule Conf-HAStates2 [THEN CompFun-ChiRel])
apply fast+
apply simp
apply (case-tac U=T)
apply simp
apply (rotate-tac -5)
apply (simp only:Source-def Target-def image-def)
apply safe
apply (rename-tac Source Trigger Guard Action Update Target)
apply (erule-tac x=Target in allE)
apply simp
apply (frule HPT-source-target2)
apply fast+
apply (rule HAStates-CompFun-SAs-mem)
apply (rule Conf-HAStates2)
apply fast+
apply (frule ChiStar-ChiPlus-noteq)
apply fast
apply (case-tac U=S)
apply (fast intro:ChiStar-Self ChiRel-ChiPlus-OneAncestor ChiPlus-ChiStar)+
done

```

lemma *SucState-StepConf*:

```

[[ TS ∈ HPT ST; S ∈ StepConf (HA ST) (Conf ST) TS;
   A ∈ the (CompFun (HA ST) S) ]] ⇒
  ∃ x. x ∈ States A ∧ x ∈ StepConf (HA ST) (Conf ST) TS
apply (case-tac S ∈ Target TS)
apply (fast intro: InitSucState-Target-StepConf InitState-States)
apply (case-tac S ∈ Conf ST)
prefer 2
apply (fast intro: InitSucState-StepConf InitState-States)
apply (case-tac S ∈ ChiStar (HA ST) “ (Source TS))
apply (fast intro: InitSucState-Conf-StepConf InitState-States)
apply (case-tac States A ∩ ChiStar (HA ST) “ (Source TS) = {})
apply (fast intro: SucState-Conf-StepConf SucState-Conf-Source-StepConf)+
done

```

8.2.18 StepStatus

lemma *StepStatus-expand*:

```

Abs-status (HA ST, StepConf (HA ST) (Conf ST) TS,

```

```

      StepActEvent TS, U !!! (Value ST))
    = (StepStatus ST TS U)
apply (unfold StepStatus-def Let-def)
apply (subst Rep-status-tuple)
apply auto
done

lemma UniqueSucState-Conf-Source-StepConf:
  [| TS ∈ HPT ST; S ∈ StepConf (HA ST) (Conf ST) TS; A ∈ SAs (HA ST);
    A ∈ the (CompFun (HA ST) S); T ∈ States A; U ∈ States A;
    T ∈ StepConf (HA ST) (Conf ST) TS; T ≠ U; U ∈ Conf ST |] ⇒
    U ∈ ChiStar (HA ST) “ Source TS
apply (frule-tac ?S2.0=T in StepConf-HAStates [THEN subsetD, THEN Comp-
  Fun-ChiRel])
apply fast+
apply (frule-tac ?S2.0=U in StepConf-HAStates [THEN subsetD, THEN Comp-
  Fun-ChiRel])
apply fast+
apply (frule-tac S=S and T=U in Conf-ChiRel, fast)
apply (case-tac S ∈ ChiStar (HA ST) “ Source TS)
apply (fast intro: ChiRel-ChiStar-trans)
apply (case-tac U ∈ Source TS)
apply force
apply (unfold StepConf-def)
apply simp
apply safe
apply (fast intro: HPT-SAStates-Target-Source)+
apply (rename-tac V)
apply (case-tac V=S)
apply (frule-tac S=S in HPT-Conf-Target-Source, fast+)
apply (fast intro: ChiStar-Image ChiRel-OneAncestor)+
apply (rename-tac V W)
apply (frule trancl-Int-mem, fold ChiPlus-def, safe)
apply (cut-tac ST=ST and S=S in HPT-Conf-Target-Source-ChiPlus)
apply fast+
apply (simp only:Image-def, safe)
apply (case-tac (V, T) ∉ ChiRel (HA ST))
apply (frule-tac S=V and T=T in ChiPlus-ChiRel-Ex)
apply (fast, safe)
apply (rename-tac X)
apply (case-tac X=S)
apply (rule-tac x=W in bexI)
prefer 4
apply (case-tac V=S)
prefer 2
apply simp
apply (fast intro: ChiPlus-ChiRel ChiRel-ChiPlus-trans2 ChiRel-OneAncestor)+
done

```

lemma *UniqueSucState-Target-StepConf*:

$\llbracket TS \in \text{HPT } ST; S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS; A \in \text{SAs } (HA \ ST);$
 $A \in \text{the } (CompFun \ (HA \ ST) \ S); T \in \text{States } A; U \in \text{States } A;$
 $T \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS; T \neq U \rrbracket \implies$
 $U \notin \text{Target } TS$

apply *auto*

apply (*frule-tac* $ST=ST$ **in** *Target-StepConf*)

apply (*subst* (*asm*) (2) *StepConf-def*)

apply *simp*

apply *safe*

apply (*cut-tac* $TS=TS$ **and** $ST=ST$ **and** $S=S$ **and** $T=U$ **in** *UniqueSucState-Conf-Source-StepConf*)

apply *fast+*

apply (*simp* *add*: *OneState-HPT-Target*)

apply (*simp* *only*: *OneState-HPT-Target-ChiRel*)

apply (*rename-tac* $V \ W$)

apply (*frule-tac* $U=W$ **and** $S=V$ **and** $T=T$ **in** *ChiRel-ChiPlus-trans2*)

apply (*frule* *trancl-Int-mem*, *fold* *ChiPlus-def*, *force*)

apply (*simp* *only*: *OneState-HPT-Target-ChiPlus*)

done

lemma *UniqueSucState-Target-ChiRel-StepConf*:

$\llbracket TS \in \text{HPT } ST; S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS; A \in \text{SAs } (HA \ ST);$
 $A \in \text{the } (CompFun \ (HA \ ST) \ S); T \in \text{States } A; U \in \text{States } A;$
 $T \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS; T \neq U; (V, U) \in \text{ChiRel } (HA \ ST);$
 $U \in \text{HAINitStates } (HA \ ST) \rrbracket$
 $\implies V \notin \text{Target } TS$

apply *auto*

apply (*frule-tac* $A=HA \ ST$ **and** $C=Conf \ ST$ **in** *Target-ChiRel-HAINit-StepConf*)

apply *fast+*

apply (*subst* (*asm*) (2) *StepConf-def*, *safe*)

apply (*fast* *intro*: *UniqueSucState-Conf-Source-StepConf*)

apply (*simp* *only*: *OneState-HPT-Target-ChiRel*)

apply (*fast* *intro*: *OneHAINitState-SAStates*)

apply (*frule* *trancl-Int-mem*, *fold* *ChiPlus-def*)

apply (*fast* *intro*: *OneHAINitState-SAStates*)

done

lemma *UniqueSucState-Target-ChiPlus-StepConf* [*rule-format*]:

$\llbracket TS \in \text{HPT } ST; (S, T) \in \text{ChiRel } (HA \ ST); (S, U) \in \text{ChiRel } (HA \ ST);$
 $V \in \text{Target } TS; (V, W) \in \text{ChiRel } (HA \ ST); T \notin \text{ChiStar } (HA \ ST) \text{ “ Source}$
 $TS;$

$(W, U) \in (\text{ChiRel } (HA \ ST) \cap \text{HAINitStates } (HA \ ST) \times \text{HAINitStates } (HA \ ST))^+;$

$T \in \text{Conf } ST \rrbracket \implies (S, U) \in \text{ChiRel } (HA \ ST) \longrightarrow T=U$

apply (*frule-tac* $S=S$ **and** $T=T$ **in** *Conf-ChiRel*)

apply *fast*

apply (*rule-tac* $a=W$ **and** $b=U$ **and** $r=\text{ChiRel } (HA \ ST) \cap \text{HAINitStates } (HA \ ST) \times \text{HAINitStates } (HA \ ST)$ **in** *trancl-induct*)

apply *safe*

```

apply (rename-tac X)
apply (case-tac W=S)
apply simp
prefer 2
apply (simp add: ChiRel-OneAncestor)
prefer 2
apply (rename-tac X Y)
apply (case-tac X=S)
apply simp
prefer 2
apply (simp add: ChiRel-OneAncestor)
prefer 2
apply (frule-tac a=V in ChiRel-HAStates)
apply (unfold HAStates-def)
apply (simp,safe)
apply (rename-tac Y)
apply (case-tac States Y  $\cap$  Source TS = {})
apply (simp add:OneState-HPT-Target-Source)
apply (subst (asm) Int-def, safe)
apply (rename-tac Z)
apply (frule-tac S=V and T=S in Conf-ChiRel)
apply fast
apply (frule HPT-Conf-Target-Source)
prefer 2
apply fast
apply fast
apply (frule-tac S=Z and T=V in OneState-HPT-SA)
apply fast+
apply simp
apply (fast intro: ChiPlus-ChiRel ChiRel-ChiPlus-trans ChiPlus-ChiStar)
apply (simp add: Image-def)
apply (frule trancl-Int-mem, fold ChiPlus-def, simp, safe)
back
apply (frule-tac T=W and S=S in Conf-ChiPlus)
apply simp
apply (frule-tac S=V and T=W in Conf-ChiRel)
apply fast
apply (frule-tac a=V in ChiRel-HAStates)
apply (unfold HAStates-def)
apply (simp, safe)
apply (rename-tac Z)
apply (case-tac States Z  $\cap$  Source TS = {})
apply (simp add:OneState-HPT-Target-Source)
apply (subst (asm) Int-def, safe)
apply (frule-tac S=V in HPT-Conf-Target-Source)
apply fast+
apply (rename-tac P)
apply (frule-tac S=P and T=V in OneState-HPT-SA)
apply fast+

```

```

apply (frule-tac  $U=V$  and  $S=W$  and  $T=S$  in ChiRel-ChiPlus-trans2)
apply fast+
apply (fast intro: ChiPlus-ChiRel ChiRel-ChiPlus-trans ChiPlus-ChiStar)
apply (case-tac  $T=S$ )
apply (simp add: ChiRel-OneAncestor)+
done

```

lemma *UniqueSucStates-SASates-StepConf*:

```

  [[  $TS \in \text{HPT } ST$ ;  $S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$ ;  $A \in \text{SAs } (HA \ ST)$ ;
     $A \in \text{the } (CompFun \ (HA \ ST) \ S)$ ;  $T \in \text{States } A$ ;  $U \in \text{States } A$ ;
     $T \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$ ;  $T \neq U$  ]]  $\implies$ 
     $U \notin \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$ 
apply (subst StepConf-def)
apply (simp, safe)
apply (rule UniqueSucState-Conf-Source-StepConf)
apply fast+
apply (frule-tac  $U=U$  in UniqueSucState-Target-StepConf)
apply fast+
apply (frule-tac  $U=U$  in UniqueSucState-Target-ChiRel-StepConf)
apply fast+
apply (rename-tac  $V \ W$ )
apply (frule trancl-Int-mem, fold ChiPlus-def)
apply (simp, safe)
apply (frule-tac  $?S2.0=T$  in StepConf-HASates [THEN subsetD, THEN Comp-Fun-ChiRel])
apply fast+
apply (frule-tac  $?S2.0=U$  in StepConf-HASates [THEN subsetD, THEN Comp-Fun-ChiRel])
apply fast+
apply (subst (asm) (2) StepConf-def)
apply (simp, safe)
apply (fast intro: UniqueSucState-Target-ChiPlus-StepConf)
apply (frule-tac  $U=W$  and  $T=U$  and  $S=T$  in OneState-HPT-Target-ChiPlus)
apply (fast intro: ChiPlus-ChiRel ChiRel-ChiPlus-trans2 OneHAINitState-SASates)+
apply (frule trancl-Int-mem, fold ChiPlus-def, simp, safe)
apply (fast intro: OneHAINitState-SASates)
done

```

lemma *UniqueSucStates-Ancestor-StepConf*:

```

  [[  $TS \in \text{HPT } ST$ ;  $S \in \text{HASates } (HA \ ST)$ ;  $SA \in \text{the } (CompFun \ (HA \ ST) \ S)$ ;
     $T \in \text{States } SA$ ;  $T \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$  ]]
   $\implies S \in \text{StepConf } (HA \ ST) \ (Conf \ ST) \ TS$ 
apply (rule notnotD, rule notI)
apply (subst (asm) StepConf-def)
apply (simp, safe)
apply (frule CompFun-ChiRel, fast+)
apply (frule Conf-ChiRel, fast)
apply (frule ChiRel-ChiStar-Source-notmem, fast+)
apply (subst (asm) StepConf-def)

```



```

apply force
apply (case-tac States SA  $\cap$  Source TS = {})
apply (simp add: OneState-HPT-Target-Source)
apply (subst (asm) Int-def)
apply (simp, safe)
apply (rename-tac U)
apply (frule-tac ?S2.0=U in CompFun-ChiRel, fast+)
apply (frule Conf-ChiRel)
apply (frule HPT-Source-Conf, fast)
apply (case-tac S  $\in$  ChiStar (HA ST) “ Source TS)
apply (subst (asm) StepConf-def)
apply simp
apply (frule ChiRel-ChiStar-notmem)
apply fast+
apply (case-tac U=S)
apply (subst (asm) StepConf-def)
apply force
apply (subst (asm) StepConf-def)
apply force
apply (rename-tac U)
apply (case-tac U=S)
apply (subst (asm) StepConf-def)
apply force
apply (frule CompFun-ChiRel, fast+)
apply (simp add: ChiRel-OneAncestor)
apply (rename-tac U V)
apply (frule trancl-Int-mem, fold ChiPlus-def, simp, safe)
apply (frule tranclD2)
apply safe
apply (rename-tac W)
apply (case-tac W=S)
apply simp
prefer 2
apply (frule CompFun-ChiRel, fast+)
apply (simp only: ChiRel-OneAncestor)
apply (subst (asm) StepConf-def)
apply safe
apply (simp add: Image-def)
apply (erule-tac x=U in ballE)
apply (case-tac U=S)
apply fast
apply (simp add: rtrancl-eq-or-trancl)
apply fast
apply (simp add: Image-def)
apply (rename-tac W)
apply (erule-tac x=U in ballE)
apply (simp add: rtrancl-eq-or-trancl)
apply fast+
done

```

```

lemma UniqueSucStates-StepConf:
  [|  $TS \in \text{HPT } ST$  |]  $\implies$ 
    UniqueSucStates (SAs (HA ST)) (CompFun (HA ST)) (StepConf (HA ST))
  (Conf ST) TS
apply (unfold UniqueSucStates-def)
apply auto
apply (simp only: SucState-StepConf)
apply (rule notnotD, rule notI)
apply (frule UniqueSucStates-SASates-StepConf)
apply fast
prefer 2
apply fast
apply (rule HASates-CompFun-SAs-mem)
prefer 2
apply fast
apply (simp only: HASates-def, fast)
apply fast+
back
apply (frule UniqueSucStates-Ancestor-StepConf)
prefer 2
apply fast
apply (simp only: HASates-def, fast)
apply fast+
done

```

```

lemma Status-Step:
  [|  $TS \in \text{HPT } ST$ ;  $U \in \text{ResolveRacing } TS$  |]  $\implies$ 
    (HA ST, StepConf (HA ST) (Conf ST) TS, StepActEvent TS, U !!! (Value
  ST))  $\in$  status
apply (unfold status-def Status-def)
apply auto
apply (frule StepActEvent-HAEvents)
apply blast
apply (unfold IsConfSet-def)
apply (rule conjI, frule StepConf-HASates, unfold HASates-def, assumption)
apply (rule conjI, rule RootExSem-StepConf, assumption)
apply (rule UniqueSucStates-StepConf, assumption)
done

```

8.3 Meta Theorem: Preservation for Statecharts

```

lemma IsConfSet-StepConf:
   $TS \in \text{HPT } ST \implies$  IsConfSet (SAs (HA ST)) (CompFun (HA ST))
    (StepConf (HA ST) (Conf ST) TS)
apply (unfold IsConfSet-def)
apply auto
apply (frule StepConf-HASates)
apply (unfold HASates-def, fast)

```

```

apply (rule RootExSem-StepConf, assumption)
apply (rule UniqueSucStates-StepConf, assumption)
done

```

```

lemma HA-StepStatus-HPT-ResolveRacing [simp]:

```

```

  [ TS ∈ HPT ST; U ∈ ResolveRacing TS ] ⇒

```

```

  HA (StepStatus ST TS U) = HA ST

```

```

apply (subst StepStatus-expand [THEN sym])

```

```

apply (subst HA-def)

```

```

apply auto

```

```

apply (subst Abs-status-inverse)

```

```

apply auto

```

```

apply (rule Status-Step)

```

```

apply auto

```

```

done

```

```

end

```

9 Kripke Structures and CTL

```

theory Kripke

```

```

imports Main

```

```

begin

```

```

definition

```

```

  Kripke :: ['s set,
            'a set,
            ('s * 's) set,
            ('s → 'a set)]
    => bool where

```

```

  Kripke S S0 R L =
    (S0 ⊆ S ∧
     R ≤ S × S ∧
     (Domain R) = S ∧
     (dom L) = S)

```

```

lemma Kripke-EmptySet:

```

```

  ({@x. True}, {@x. True}, {@x. True, @x. True}, Map.empty(@x. True ↦ {@x.
  True})) ∈

```

```

  {(S,S0,R,L) | S S0 R L. Kripke S S0 R L}

```

```

by (unfold Kripke-def Domain-unfold, auto)

```

```

definition

```

```

  kripke =
    {(S,S0,T,L) |
     (S::('s set))
     (S0::('s set))
     (T::(('s * 's) set))

```

($L::('s \rightarrow ('a \text{ set}))$).
Kripke S S0 T L}

typedef ('s,'a) *kripke* =
kripke :: ('s set * 's set * ('s * 's) set * ('s \rightarrow 'a set)) set
unfolding *kripke-def*
apply (rule *exI*)
apply (rule *Kripke-EmptySet*)
done

definition
Statuses :: ('s,'a) *kripke* \Rightarrow 's set **where**
Statuses = *fst o Rep-kripke*

definition
InitStatuses :: ('s,'a) *kripke* \Rightarrow 's set **where**
InitStatuses == *fst o snd o Rep-kripke*

definition
StepRel :: ('s,'a) *kripke* \Rightarrow ('s * 's) set **where**
StepRel == *fst o snd o snd o Rep-kripke*

definition
LabelFun :: ('s,'a) *kripke* \Rightarrow ('s \rightarrow 'a set) **where**
LabelFun == *snd o snd o snd o Rep-kripke*

definition
Paths :: [('s,'a) *kripke*, 's] \Rightarrow
(nat \Rightarrow 's) set **where**
Paths M S == { *p* . *S* = *p* (0::nat) \wedge (\forall *i*. (*p* *i*, *p* (*i*+1)) \in (*StepRel M*)) }

datatype ('s,'a) *ctl* = *Atom* 'a
| *AND* ('s,'a) *ctl* ('s,'a) *ctl*
| *OR* ('s,'a) *ctl* ('s,'a) *ctl*
| *IMPLIES* ('s,'a) *ctl* ('s,'a) *ctl*
| *CAX* ('s,'a) *ctl*
| *AF* ('s,'a) *ctl*
| *AG* ('s,'a) *ctl*
| *AU* ('s,'a) *ctl* ('s,'a) *ctl*
| *AR* ('s,'a) *ctl* ('s,'a) *ctl*

primrec
eval-ctl :: [('s,'a) *kripke*, 's, ('s,'a) *ctl*] \Rightarrow bool (-,- |=c= - [92,91,90]90)
where
(*M,S* |=c= (*Atom P*)) = (*P* \in the ((*LabelFun M*) *S*))
| (*M,S* |=c= (*AND F1 F2*)) = ((*M,S* |=c= *F1*) \wedge (*M,S* |=c= *F2*))
| (*M,S* |=c= (*OR F1 F2*)) = ((*M,S* |=c= *F1*) \vee (*M,S* |=c= *F2*))
| (*M,S* |=c= (*IMPLIES F1 F2*)) = ((*M,S* |=c= *F1*) \longrightarrow (*M,S* |=c= *F2*))

$$\begin{aligned}
& | (M, S \models_{c=} (CAX F)) & = (\forall T. (S, T) \in (StepRel M) \longrightarrow (M, T \models_{c=} \\
& F)) \\
& | (M, S \models_{c=} (AF F)) & = (\forall P \in Paths M S. \exists i. (M, (P i) \models_{c=} F)) \\
& | (M, S \models_{c=} (AG F)) & = (\forall P \in Paths M S. \forall i. (M, (P i) \models_{c=} F)) \\
& | (M, S \models_{c=} (AU F G)) & = (\forall P \in Paths M S. \\
& & \exists i. (M, (P i) \models_{c=} G) \wedge \\
& & (\forall j. j < i \longrightarrow (M, (P j) \models_{c=} F))) \\
& | (M, S \models_{c=} (AR F G)) & = (\forall P \in Paths M S. \\
& & \forall i. (M, (P i) \models_{c=} G) \vee \\
& & (\exists j. j < i \wedge (M, (P j) \models_{c=} F)))
\end{aligned}$$

end

10 Kripke Structures as Hierarchical Automata

theory *HAKripke*

imports *HASem Kripke*

begin

type-synonym $(s, e, d)hakripke = ((s, e, d)status, (s, e, d)atomar)kripke$

type-synonym $(s, e, d)hactl = ((s, e, d)status, (s, e, d)atomar)ctl$

definition

LabelFunSem :: $(s, e, d)hierauto$
 $\Rightarrow ((s, e, d)status \rightarrow (((s, e, d) atomar) set))$ **where**
LabelFunSem a = $(\lambda ST.$
 (if (HA ST = a) then
 (let
 In-preds = $(\lambda s. (IN s)) \text{ ' } (Conf ST);$
 En-preds = $(\lambda e. (EN e)) \text{ ' } (Events ST);$
 Val-preds = $\{ x . (\exists P. (x = (VAL P)) \wedge P (Value ST)) \}$
 in
 Some (In-preds \cup En-preds \cup Val-preds \cup {atomar.TRUE})
 else
 None))

definition

HA2Kripke :: $(s, e, d)hierauto \Rightarrow (s, e, d)hakripke$ **where**
HA2Kripke a =
 Abs-kripke ($\{ST. HA ST = a\},$
 {InitStatus a},
 StepRelSem a,
 LabelFunSem a)

definition

eval-ctl-HA :: $[(s, e, d)hierauto, (s, e, d)hactl]$
 $\Rightarrow bool (- \models_H = - [92, 91] 90)$ **where**

eval-ctl-HA a f = $((HA2Kripke a), (InitStatus a) \models_{c=} f)$

```

lemma Kripke-HA [simp]:
  Kripke {ST. HA ST = a} {InitStatus a} (StepRelSem a) (LabelFunSem a)
apply (unfold Kripke-def)
apply auto
apply (unfold StepRelSem-def)
apply auto
apply (unfold LabelFunSem-def Let-def If-def dom-def)
apply auto
prefer 2
apply (rename-tac ST S)
apply (case-tac HA ST = a)
apply auto
apply (rename-tac ST)
apply (case-tac HPT ST = {})
apply auto
apply (rename-tac TSS)
apply (erule-tac x=StepStatus ST TSS (@ u. u : ResolveRacing TSS) in allE)
apply (erule-tac x=TSS in ballE)
apply auto
done

```

```

lemma LabelFun-LabelFunSem [simp]:
  (LabelFun (HA2Kripke a)) = (LabelFunSem a)
apply (unfold HA2Kripke-def LabelFun-def)
apply auto
apply (subst Abs-kripke-inverse)
apply auto
apply (unfold kripke-def)
apply auto
done

```

```

lemma InitStatuses-InitStatus [simp]:
  (InitStatuses (HA2Kripke a)) = {(InitStatus a)}
apply (unfold HA2Kripke-def InitStatuses-def)
apply simp
apply (subst Abs-kripke-inverse)
apply (unfold kripke-def)
apply auto
done

```

```

lemma Statuses-StatusesOfHA [simp]:
  (Statuses (HA2Kripke a)) = {ST. HA ST = a}
apply (unfold HA2Kripke-def Statuses-def)
apply simp
apply (subst Abs-kripke-inverse)
apply (unfold kripke-def)
apply auto
done

```

lemma *StepRel-StepRelSem* [simp]:
 $(\text{StepRel } (\text{HA2Kripke } a)) = (\text{StepRelSem } a)$
apply (unfold HA2Kripke-def StepRel-def)
apply simp
apply (subst Abs-kripke-inverse)
apply (unfold kripke-def)
apply auto
done

lemma *TRUE-LabelFunSem* [simp]:
 $\text{atomar.TRUE} \in \text{the } (\text{LabelFunSem } (\text{HA } ST) ST)$
apply (unfold LabelFunSem-def Let-def)
apply auto
done

lemma *FALSE-LabelFunSem* [simp]:
 $\text{atomar.FALSE} \notin \text{the } (\text{LabelFunSem } (\text{HA } ST) ST)$
apply (unfold LabelFunSem-def Let-def)
apply auto
done

lemma *Conf-LabelFunSem* [simp]:
 $((\text{IN } S) \in \text{the } (\text{LabelFunSem } (\text{HA } ST) ST)) = (S \in (\text{Conf } ST))$
apply (unfold LabelFunSem-def Let-def)
apply auto
done

lemma *Events-LabelFunSem* [simp]:
 $((\text{EN } S) \in \text{the } (\text{LabelFunSem } (\text{HA } ST) ST)) = (S \in (\text{Events } ST))$
apply (unfold LabelFunSem-def Let-def)
apply auto
done

lemma *Value-LabelFunSem* [simp]:
 $((\text{VAL } P) \in \text{the } (\text{LabelFunSem } (\text{HA } ST) ST)) = (P \in (\text{Value } ST))$
apply (unfold LabelFunSem-def Let-def)
apply auto
done

lemma *AtomTRUE-EvalCTLHA* [simp]:
 $a \models_H (\text{Atom } (\text{atomar.TRUE}))$
apply (unfold eval-ctl-HA-def)
apply auto
apply (subst HA-InitStatus [THEN sym])
apply (rule TRUE-LabelFunSem)
done

lemma *AtomFalse-EvalCTLHA* [simp]:

```

  ¬ a |=H= (Atom (atomar.FALSE))
apply (unfold eval-ctl-HA-def)
apply auto
apply (subst (asm) HA-InitStatus [THEN sym])
apply (simp only: FALSE-LabelFunSem)
done

```

```

lemma Events-InitStatus-EvalCTLHA [simp]:
  (a |=H= (Atom (EN S))) = (S ∈ (Events (InitStatus a)))
apply (unfold eval-ctl-HA-def)
apply simp
apply (subst HA-InitStatus [THEN sym])
apply (rule Events-LabelFunSem)
done

```

```

lemma Conf-InitStatus-EvalCTLHA [simp]:
  (a |=H= (Atom (IN S))) = (S ∈ (Conf (InitStatus a)))
apply (unfold eval-ctl-HA-def)
apply simp
apply (subst HA-InitStatus [THEN sym])
apply (subst Conf-LabelFunSem)
apply simp
done

```

```

lemma HAINitValue-EvalCTLHA [simp]:
  (a |=H= (Atom (VAL P))) = (P (HAINitValue a))
apply (unfold eval-ctl-HA-def)
apply simp
apply (subst HA-InitStatus [THEN sym])
apply (subst Value-LabelFunSem)
apply auto
done

```

end

11 Constructing Hierarchical Automata

```

theory HAOps
imports HA
begin

```

11.1 Constructing a Composition Function for a PseudoHA

definition

```

  EmptyMap :: 's set => ('s → (('s,'e,'d)seqauto) set) where
  EmptyMap S = (λ a . if a ∈ S then Some {} else None)

```

```

lemma EmptyMap-dom [simp]:
  dom (EmptyMap S) = S

```


by (unfold dom-def EmptyMap-def, auto)

lemma *EmptyMap-ran* [simp]:

$S \neq \{\} \implies \text{ran } (\text{EmptyMap } S) = \{\{\}\}$

by (unfold ran-def EmptyMap-def, auto)

lemma *EmptyMap-the* [simp]:

$x \in S \implies \text{the } ((\text{EmptyMap } S) x) = \{\}$

by (unfold ran-def EmptyMap-def, auto)

lemma *EmptyMap-ran-override*:

$\llbracket S \neq \{\}; (S \cap (\text{dom } G)) = \{\} \rrbracket \implies$

$\text{ran } (G ++ \text{EmptyMap } S) = \text{insert } \{\} (\text{ran } G)$

apply (subst ran-override)

apply (simp add: Int-commute)

apply simp

done

lemma *EmptyMap-Union-ran-override*:

$\llbracket S \neq \{\};$

$S \cap \text{dom } G = \{\} \rrbracket \implies$

$(\text{Union } (\text{ran } (G ++ (\text{EmptyMap } S)))) = (\text{Union } (\text{ran } G))$

apply (subst EmptyMap-ran-override)

apply auto

done

lemma *EmptyMap-Union-ran-override2*:

$\llbracket S \neq \{\}; S \cap \text{dom } G1 = \{\};$

$\text{dom } G1 \cap \text{dom } G2 = \{\} \rrbracket \implies$

$\bigcup (\text{ran } (G1 ++ \text{EmptyMap } S ++ G2)) = (\bigcup (\text{ran } G1 \cup \text{ran } G2))$

apply (unfold Union-eq UNION-eq EmptyMap-def Int-def ran-def)

apply (simp add: map-add-Some-iff)

apply (unfold dom-def)

apply simp

apply (rule equalityI)

apply (rule subsetI)

apply simp

apply fast

apply (rule subsetI)

apply (rename-tac t)

apply simp

apply (erule bexE)

apply (rename-tac U)

apply simp

apply (erule disjE)

apply (erule exE)

apply (rename-tac v)

apply (rule-tac $x=U$ in exI)

apply simp

apply (*rule-tac* $x=v$ **in** exI)
apply *auto*
done

lemma *EmptyMap-Root* [*simp*]:
 $Root \{SA\} (EmptyMap (States SA)) = SA$
by (*unfold Root-def, auto*)

lemma *EmptyMap-RootEx* [*simp*]:
 $RootEx \{SA\} (EmptyMap (States SA))$
by (*unfold RootEx-def, auto*)

lemma *EmptyMap-OneAncestor* [*simp*]:
 $OneAncestor \{SA\} (EmptyMap (States SA))$
by (*unfold OneAncestor-def, auto*)

lemma *EmptyMap-NoCycles* [*simp*]:
 $NoCycles \{SA\} (EmptyMap (States SA))$
by (*unfold NoCycles-def EmptyMap-def, auto*)

lemma *EmptyMap-IsCompFun* [*simp*]:
 $IsCompFun \{SA\} (EmptyMap (States SA))$
by (*unfold IsCompFun-def, auto*)

lemma *EmptyMap-hierauto* [*simp*]:
 $(D, \{SA\}, SAEvents SA, EmptyMap (States SA)) \in hierauto$
by (*unfold hierauto-def HierAuto-def, auto*)

11.2 Extending a Composition Function by a SA

definition

$$\begin{aligned}
 FAddSA &:: [(s \rightarrow ((s, 'e, 'd)seqauto) set), 's * (s, 'e, 'd)seqauto] \\
 &=> (s \rightarrow ((s, 'e, 'd)seqauto) set) \\
 &((\text{- } [f+] / \text{-}) [10, 11] 10) \textbf{ where} \\
 FAddSA \ G \ SSA &= (\text{let } (S, SA) = SSA \\
 &\quad \text{in} \\
 &\quad \text{if } ((S \in \text{dom } G) \wedge (S \notin \text{States } SA)) \text{ then} \\
 &\quad \quad (G ++ (\text{Map.empty}(S \mapsto (\text{insert } SA (\text{the } (G \ S)))))) \\
 &\quad \quad ++ \text{EmptyMap } (\text{States } SA)) \\
 &\quad \text{else } G)
 \end{aligned}$$

lemma *FAddSA-dom* [*simp*]:
 $(S \notin (\text{dom } (A::(a => (a, 'c, 'd)seqauto) set option)))) \implies$
 $((A [f+] (S, (SA::(a, 'c, 'd)seqauto))) = A)$
by (*unfold FAddSA-def Let-def, auto*)

lemma *FAddSA-States* [*simp*]:
 $(S \in (\text{States } (SA::(a, 'c, 'd)seqauto))) \implies$
 $((A::(a => (a, 'c, 'd)seqauto) set option) [f+] (S, SA) = A)$

by (*unfold FAddSA-def Let-def, auto*)

lemma *FAddSA-dom-insert [simp]*:

$\llbracket S \in (\text{dom } A); S \notin \text{States } SA \rrbracket \implies$
 $((A [f+] (S, SA)) S) = \text{Some } (\text{insert } SA (\text{the } (A S)))$
by (*unfold FAddSA-def Let-def restrict-def, auto*)

lemma *FAddSA-States-neq [simp]*:

$\llbracket S' \notin \text{States } (SA::('a, 'c, 'd)\text{seqauto}); S \neq S' \rrbracket \implies$
 $((A::('a \Rightarrow ('a, 'c, 'd)\text{seqauto set option})) [f+] (S, SA)) S') = (A S')$
apply (*case-tac S \in dom A*)
apply (*case-tac S \in States SA*)
apply *auto*
apply (*case-tac S' \in dom A*)
apply (*unfold FAddSA-def Let-def*)
apply *auto*
apply (*simp add: dom-None*)
done

lemma *FAddSA-dom-emptyset [simp]*:

$\llbracket S \in (\text{dom } A); S \notin \text{States } SA; S' \in \text{States } (SA::('a, 'c, 'd)\text{seqauto}) \rrbracket \implies$
 $((A::('a \Rightarrow ('a, 'c, 'd)\text{seqauto set option})) [f+] (S, SA)) S') = (\text{Some } \{\})$
apply (*unfold FAddSA-def Let-def*)
apply *auto*
apply (*unfold EmptyMap-def*)
apply *auto*
done

lemma *FAddSA-dom-dom-States [simp]*:

$\llbracket S \in (\text{dom } F); S \notin \text{States } SA \rrbracket \implies$
 $(\text{dom } ((F::('a \rightarrow (('a, 'b, 'd)\text{seqauto set})) [f+] (S, SA))) =$
 $((\text{dom } F) \cup (\text{States } (SA::('a, 'b, 'd)\text{seqauto}))))$
by (*unfold FAddSA-def Let-def, auto*)

lemma *FAddSA-dom-dom [simp]*:

$S \notin (\text{dom } F) \implies$
 $(\text{dom } ((F::('a \rightarrow (('a, 'b, 'd)\text{seqauto set})) [f+] (S, (SA::('a, 'b, 'd)\text{seqauto})))) = (\text{dom } F)$
by (*unfold FAddSA-def Let-def, auto*)

lemma *FAddSA-States-dom [simp]*:

$S \in (\text{States } SA) \implies$
 $(\text{dom } ((F::('a \rightarrow (('a, 'b, 'd)\text{seqauto set})) [f+] (S, (SA::('a, 'b, 'd)\text{seqauto})))) = (\text{dom } F)$
by (*unfold FAddSA-def Let-def, auto*)

lemma *FAddSA-dom-insert-dom-disjunct [simp]*:

$\llbracket S \in \text{dom } G; \text{States } SA \cap \text{dom } G = \{\} \rrbracket \implies ((G [f+] (S, SA)) S) = \text{Some } (\text{insert } SA (\text{the } (G S)))$

```

apply (rule FAddSA-dom-insert)
apply auto
done

```

```

lemma FAddSA-Union-ran:
   $\llbracket S \in \text{dom } G; (\text{States } SA) \cap (\text{dom } G) = \{\} \rrbracket \implies$ 
   $(\bigcup (\text{ran } (G [f+] (S,SA)))) = (\text{insert } SA (\bigcup (\text{ran } G)))$ 
apply (unfold FAddSA-def Let-def)
apply simp
apply (rule conjI)
prefer 2
apply (rule impI)
apply (unfold Int-def)
apply simp
apply (fold Int-def)
apply (rule impI)
apply (subst EmptyMap-Union-ran-override)
apply auto
done

```

```

lemma FAddSA-Union-ran2:
   $\llbracket S \in \text{dom } G1; (\text{States } SA) \cap (\text{dom } G1) = \{\}; (\text{dom } G1 \cap \text{dom } G2) = \{\} \rrbracket \implies$ 
   $(\bigcup (\text{ran } ((G1 [f+] (S,SA)) ++ G2))) = (\text{insert } SA (\bigcup ((\text{ran } G1) \cup (\text{ran } G2))))$ 
apply (unfold FAddSA-def Let-def)
apply (simp (no-asm-simp))
apply (rule conjI)
apply (rule impI)
apply (subst EmptyMap-Union-ran-override2)
apply simp
apply simp
apply simp
apply fast
apply (subst Union-Un-distrib)
apply (subst Union-ran-override2)
apply auto
done

```

```

lemma FAddSA-ran:
   $\llbracket \forall T \in \text{dom } G . T \neq S \longrightarrow (\text{the } (G T) \cap \text{the } (G S)) = \{\};$ 
   $S \in \text{dom } G; (\text{States } SA) \cap (\text{dom } G) = \{\} \rrbracket \implies$ 
   $\text{ran } (G [f+] (S,SA)) = \text{insert } \{\} (\text{insert } (\text{insert } SA (\text{the } (G S))) (\text{ran } G - \{\text{the } (G S)\}))$ 
apply (unfold FAddSA-def Let-def)
apply simp
apply (rule conjI)
apply (rule impI)+
prefer 2
apply fast
apply (simp add: EmptyMap-ran-override)

```

apply (*unfold ran-def*)
apply *auto*
apply (*rename-tac Y X a xa xb*)
apply (*erule-tac x=a in allE*)
apply *simp*
apply (*erule-tac x=a in allE*)
apply *simp*
done

lemma *FAddSA-RootEx-def*:

$$\llbracket S \in \text{dom } G; (\text{States } SA) \cap (\text{dom } G) = \{\} \rrbracket \implies$$

$$\text{RootEx } F (G [f+] (S, SA)) = (\exists! A . A \in F \wedge A \notin \text{insert } SA (\bigcup (\text{ran } G)))$$
apply (*unfold RootEx-def*)
apply (*simp only: FAddSA-Union-ran Int-commute*)
done

lemma *FAddSA-RootEx*:

$$\llbracket \bigcup (\text{ran } G) = F - \{\text{Root } F G\};$$

$$\text{dom } G = \bigcup (\text{States } ' F);$$

$$(\text{dom } G \cap \text{States } SA) = \{\}; S \in \text{dom } G;$$

$$\text{RootEx } F G \rrbracket \implies \text{RootEx } (\text{insert } SA F) (G [f+] (S, SA))$$
apply (*simp add: FAddSA-RootEx-def Int-commute cong: rev-conj-cong*)
apply (*auto cong: conj-cong*)
done

lemma *FAddSA-Root-def*:

$$\llbracket S \in \text{dom } G; (\text{States } SA) \cap (\text{dom } G) = \{\} \rrbracket \implies$$

$$(\text{Root } F (G [f+] (S, SA)) = (@ A . A \in F \wedge A \notin \text{insert } SA (\bigcup (\text{ran } G))))$$
apply (*unfold Root-def*)
apply (*simp only: FAddSA-Union-ran Int-commute*)
done

lemma *FAddSA-RootEx-Root*:

$$\llbracket \text{Union } (\text{ran } G) = F - \{\text{Root } F G\};$$

$$\bigcup (\text{States } ' F) = \text{dom } G;$$

$$(\text{dom } G \cap \text{States } SA) = \{\}; S \in \text{dom } G;$$

$$\text{RootEx } F G \rrbracket \implies (\text{Root } (\text{insert } SA F) (G [f+] (S, SA))) = (\text{Root } F G)$$
apply (*simp add: FAddSA-Root-def Int-commute cong: rev-conj-cong*)
apply (*simp cong: conj-cong*)
done

lemma *FAddSA-OneAncestor*:

$$\llbracket \bigcup (\text{ran } G) = F - \{\text{Root } F G\};$$

$$(\text{dom } G \cap \text{States } SA) = \{\}; S \in \text{dom } G;$$

$$\bigcup (\text{States } ' F) = \text{dom } G; \text{RootEx } F G;$$

$$\text{OneAncestor } F G \rrbracket \implies \text{OneAncestor } (\text{insert } SA F) (G [f+] (S, SA))$$
apply (*subst OneAncestor-def*)
apply *simp*
apply (*rule ballI*)

```

apply (rename-tac SAA)
apply (case-tac SA = SAA)
apply (rule-tac a=S in ex1I)
apply (rule conjI)
apply simp
apply fast
apply (subst FAddSA-dom-insert)
apply simp
apply (simp add:Int-def)
apply simp
apply (rename-tac T)
apply (erule conjE bexE exE disjE)+
apply (rename-tac SAAA)
apply simp
apply (erule conjE)
apply (subst not-not [THEN sym])
apply (rule notI)
apply (case-tac T ∈ States SAA)
apply blast
apply (drule-tac A=G and S=S and SA=SAA in FAddSA-States-neq)
apply fast
apply simp
apply (case-tac SAA ∉ Union (ran G))
apply (frule ran-dom-the)
prefer 2
apply fast
apply blast
apply simp
apply (erule conjE)
apply (simp add: States-Int-not-mem)
apply (unfold OneAncestor-def)
apply (drule-tac G=G and S=S and SA=SA in FAddSA-RootEx-Root)
apply simp
apply simp
apply simp
apply simp
apply (erule-tac x=SAA in balle)
prefer 2
apply simp
apply simp
apply (erule conjE bexE ex1E exE disjE)+
apply (rename-tac T SAAA)
apply (rule-tac a=T in ex1I)
apply (rule conjI)
apply fast
apply (case-tac T = S)
apply simp
apply (case-tac S ∉ States SA)
apply simp

```

```

apply simp
apply (subst FAddSA-States-neq)
apply blast
apply (rule not-sym)
apply simp
apply simp
apply (rename-tac U)
apply simp
apply (erule conjE bexE)+
apply (rename-tac SAAAA)
apply simp
apply (erule conjE disjE)+
apply (frule FAddSA-dom-emptyset)
prefer 2
apply fast
back
back
apply simp
apply blast
apply simp
apply (erule-tac x=U in allE)
apply (erule impE)
prefer 2
apply simp
apply (rule conjI)
apply fast
apply (case-tac S ≠ U)
apply (subgoal-tac U ∉ States SA)
apply (drule-tac A=G in FAddSA-States-neq)
apply fast
apply simp
apply blast
apply (drule-tac A=G and SA=SA in FAddSA-dom-insert)
apply simp
apply blast
apply auto
done

```

```

lemma FAddSA-NoCycles:
  [| (States SA ∩ dom G) = {}; S ∈ dom G;
    dom G = ∪ (States ' F); NoCycles F G |] ⇒
    NoCycles (insert SA F) (G [f+] (S,SA))
apply (unfold NoCycles-def)
apply (rule ballI impI)+
apply (rename-tac SAA)
apply (case-tac ∃ s ∈ SAA. s ∈ States SA)
apply simp
apply (erule bexE)+
apply (rename-tac SAAA T)

```

```

apply (rule-tac x=T in beXI)
apply simp
apply (subst FAddSA-dom-emptyset)
apply simp
apply fast
apply blast
apply simp
apply simp
apply simp
apply simp
apply (erule-tac x=SAA in ballE)
prefer 2
apply simp
apply auto[1]
apply (unfold UNION-eq Pow-def)
apply simp
apply (case-tac SAA = {})
apply fast
apply simp
apply (erule bexE)+
apply (rename-tac SAAA T)
apply (rule-tac x=T in beXI)
prefer 2
apply simp
apply (case-tac T=S)
apply simp
apply (subst FAddSA-dom-insert)
apply auto
done

```

lemma FAddSA-IsCompFun:

$\llbracket (States SA \cap (\bigcup (States ' F))) = \{\};$

$S \in (\bigcup (States ' F));$

$IsCompFun F G \rrbracket \implies IsCompFun (insert SA F) (G [f+] (S,SA))$

```

apply (unfold IsCompFun-def)

```

```

apply (erule conjE)+

```

```

apply (simp add: Int-commute FAddSA-RootEx-Root FAddSA-RootEx FAddSA-OneAncestor
FAddSA-NoCycles)

```

```

apply (rule conjI)

```

```

apply (subst FAddSA-dom-dom-States)

```

```

apply simp

```

```

apply blast

```

```

apply (simp add: Un-commute)

```

```

apply (simp add: FAddSA-Union-ran)

```

```

apply (case-tac SA = Root F G)

```

```

prefer 2

```

```

apply blast

```

```

apply (subgoal-tac States (Root F G)  $\subseteq \bigcup (States ' F)$ )

```

```

apply simp

```


apply (*frule subset-lemma*)
apply *auto*
done

lemma *FAddSA-HierAuto*:
 $\llbracket (States\ SA \cap (\bigcup (States\ 'F))) = \{\};$
 $S \in (\bigcup (States\ 'F));$
 $HierAuto\ D\ F\ E\ G \rrbracket \implies HierAuto\ D\ (insert\ SA\ F)\ (E \cup SAEvents\ SA)\ (G$
 $[f+]\ (S,SA))$
apply (*unfold HierAuto-def*)
apply *auto*
apply (*simp add: MutuallyDistinct-Insert*)
apply (*rule FAddSA-IsCompFun*)
apply *auto*
done

lemma *FAddSA-HierAuto-insert [simp]*:
 $\llbracket (States\ SA \cap HStates\ HA) = \{\};$
 $S \in HStates\ HA \rrbracket \implies$
 $HierAuto\ (HInitValue\ HA)$
 $(insert\ SA\ (SAs\ HA))$
 $(HAEvents\ HA \cup SAEvents\ SA)$
 $(CompFun\ HA\ [f+]\ (S,SA))$
apply (*unfold HStates-def*)
apply (*rule FAddSA-HierAuto*)
apply *auto*
done

11.3 Constructing a PseudoHA

definition

PseudoHA :: $[(\ 's, 'e, 'd) seqauto, 'd\ data] \implies (\ 's, 'e, 'd) hierauto$ **where**
 $PseudoHA\ SA\ D = Abs-hierauto(D, \{SA\}, SAEvents\ SA, EmptyMap\ (States\ SA))$

lemma *PseudoHA-SAs [simp]*:
 $SAs\ (PseudoHA\ SA\ D) = \{SA\}$
by (*unfold PseudoHA-def SAs-def, simp add: Abs-hierauto-inverse*)

lemma *PseudoHA-Events [simp]*:
 $HAEvents\ (PseudoHA\ SA\ D) = SAEvents\ SA$
by (*unfold PseudoHA-def HAEvents-def, simp add: Abs-hierauto-inverse*)

lemma *PseudoHA-CompFun [simp]*:
 $CompFun\ (PseudoHA\ SA\ D) = EmptyMap\ (States\ SA)$
by (*unfold PseudoHA-def CompFun-def, simp add: Abs-hierauto-inverse*)

lemma *PseudoHA-HStates [simp]*:
 $HStates\ (PseudoHA\ SA\ D) = (States\ SA)$
by (*unfold HStates-def, auto*)

lemma *PseudoHA-HAInitValue* [simp]:
 $(HAInitValue (PseudoHA SA D)) = D$
by (*unfold PseudoHA-def Let-def HAInitValue-def, simp add: Abs-hierauto-inverse*)

lemma *PseudoHA-CompFun-the* [simp]:
 $S \in States A \implies (the (CompFun (PseudoHA A D) S)) = \{\}$
by *simp*

lemma *PseudoHA-CompFun-ran* [simp]:
 $(ran (CompFun (PseudoHA SA D))) = \{\{\}\}$
by *auto*

lemma *PseudoHA-HARoot* [simp]:
 $(HARoot (PseudoHA SA D)) = SA$
by (*unfold HARoot-def, auto*)

lemma *PseudoHA-HAInitState* [simp]:
 $HAInitState (PseudoHA A D) = InitState A$
apply (*unfold HAInitState-def*)
apply *simp*
done

lemma *PseudoHA-HAInitStates* [simp]:
 $HAInitStates (PseudoHA A D) = \{InitState A\}$
apply (*unfold HAInitStates-def*)
apply *simp*
done

lemma *PseudoHA-Chi* [simp]:
 $S \in States A \implies Chi (PseudoHA A D) S = \{\}$
apply (*unfold Chi-def restrict-def*)
apply *auto*
done

lemma *PseudoHA-ChiRel* [simp]:
 $ChiRel (PseudoHA A D) = \{\}$
apply (*unfold ChiRel-def*)
apply *simp*
done

lemma *PseudoHA-InitConf* [simp]:
 $InitConf (PseudoHA A D) = \{InitState A\}$
apply (*unfold InitConf-def*)
apply *simp*
done

11.4 Extending a HA by a SA (*AddSA*)

definition

```

AddSA :: [(s,e,d)hierauto, 's * (s,e,d)seqauto]
      => (s,e,d)hierauto
      ((- [++]/ -) [10,11]10) where
AddSA HA SSA = (let (S,SA) = SSA;
                    DNew = HAINitValue HA;
                    FNew = insert SA (SAs HA);
                    ENew = HAEvents HA ∪ SAEvents SA;
                    GNew = CompFun HA [f+] (S,SA)
                    in
                    Abs-hierauto(DNew,FNew,ENew,GNew))

```

definition

```

AddHA :: [(s,e,d)hierauto, 's * (s,e,d)hierauto]
      => (s,e,d)hierauto
      ((- [**]/ -) [10,11]10) where
AddHA HA1 SHA =
  (let (S,HA2) = SHA;
      (D1,F1,E1,G1) = Rep-hierauto (HA1 [++] (S,HARoot HA2));
      (D2,F2,E2,G2) = Rep-hierauto HA2;
      FNew = F1 ∪ F2;
      ENew = E1 ∪ E2;
      GNew = G1 ++ G2
      in
      Abs-hierauto(D1,FNew,ENew,GNew))

```

lemma AddSA-SAs:

```

[[ (States SA ∩ HAStates HA) = {}];
 S ∈ HAStates HA ] => (SAs (HA [++] (S,SA))) = insert SA (SAs HA)
apply (unfold Let-def AddSA-def)
apply (subst SAs-def)
apply (simp add: hierauto-def Abs-hierauto-inverse)
done

```

lemma AddSA-Events:

```

[[ (States SA ∩ HAStates HA) = {}];
 S ∈ HAStates HA ] =>
  HAEvents (HA [++] (S,SA)) = (HAEvents HA) ∪ (SAEvents SA)
apply (unfold Let-def AddSA-def)
apply (subst HAEvents-def)
apply (simp add: hierauto-def Abs-hierauto-inverse)
done

```

lemma AddSA-CompFun:

```

[[ (States SA ∩ HAStates HA) = {}];
 S ∈ HAStates HA ] =>
  CompFun (HA [++] (S,SA)) = (CompFun HA [f+] (S,SA))
apply (unfold Let-def AddSA-def)

```

apply (*subst CompFun-def*)
apply (*simp add: hierauto-def Abs-hierauto-inverse*)
done

lemma *AddSA-HAStates*:

$$\llbracket (States SA \cap HAStates HA) = \{\};$$

$$S \in HAStates HA \rrbracket \implies$$

$$HAStates (HA [++] (S,SA)) = (HAStates HA) \cup (States SA)$$
apply (*unfold HAStates-def*)
apply (*subst AddSA-SAs*)
apply (*unfold HAStates-def*)
apply *auto*
done

lemma *AddSA-HAInitValue*:

$$\llbracket (States SA \cap HAStates HA) = \{\};$$

$$S \in HAStates HA \rrbracket \implies$$

$$(HAInitValue (HA [++] (S,SA))) = (HAInitValue HA)$$
apply (*unfold Let-def AddSA-def*)
apply (*subst HAINitValue-def*)
apply (*simp add: hierauto-def Abs-hierauto-inverse*)
done

lemma *AddSA-HARoot*:

$$\llbracket (States SA \cap HAStates HA) = \{\};$$

$$S \in HAStates HA \rrbracket \implies$$

$$(HARoot (HA [++] (S,SA))) = (HARoot HA)$$
apply (*unfold HARoot-def*)
apply (*simp add: AddSA-CompFun AddSA-SAs*)
apply (*subst FAddSA-RootEx-Root*)
apply *auto*
apply (*simp only: HAStates-SA-mem*)
apply (*unfold HAStates-def*)
apply *fast*
done

lemma *AddSA-CompFun-the*:

$$\llbracket (States SA \cap HAStates A) = \{\};$$

$$S \in HAStates A \rrbracket \implies$$

$$(the ((CompFun (A [++] (S,SA))) S)) = insert SA (the ((CompFun A) S))$$
by (*simp add: AddSA-CompFun*)

lemma *AddSA-CompFun-the2*:

$$\llbracket S' \in States (SA::('a,'c,'d)seqauto);$$

$$(States SA \cap HAStates A) = \{\};$$

$$S \in HAStates A \rrbracket \implies$$

$$the ((CompFun (A [++] (S,SA))) S') = \{\}$$
apply (*simp add: AddSA-CompFun*)
apply (*subst FAddSA-dom-emptyset*)

apply *auto*
done

lemma *AddSA-CompFun-the3:*

$\llbracket S' \notin \text{States } (SA::('a,'c,'d)\text{seqauto});$
 $S \neq S';$
 $(\text{States } SA \cap \text{HAStates } A) = \{\};$
 $S \in \text{HAStates } A \rrbracket \implies$
 $(\text{the } ((\text{CompFun } (A \text{ [++]} (S,SA))) S')) = (\text{the } ((\text{CompFun } A) S'))$
by (*simp add: AddSA-CompFun*)

lemma *AddSA-CompFun-ran:*

$\llbracket (\text{States } SA \cap \text{HAStates } A) = \{\};$
 $S \in \text{HAStates } A \rrbracket \implies$
 $\text{ran } (\text{CompFun } (A \text{ [++]} (S,SA))) =$
 $\text{insert } \{\} (\text{insert } (\text{insert } SA (\text{the } ((\text{CompFun } A) S))) (\text{ran } (\text{CompFun } A) -$
 $\{\text{the } ((\text{CompFun } A) S)\}))$
apply (*simp add: AddSA-CompFun*)
apply (*subst FAddSA-ran*)
apply *auto*
apply (*fast dest: CompFun-Int-disjoint*)
done

lemma *AddSA-CompFun-ran2:*

$\llbracket (\text{States } SA1 \cap \text{HAStates } A) = \{\};$
 $(\text{States } SA2 \cap (\text{HAStates } A \cup \text{States } SA1)) = \{\};$
 $S \in \text{HAStates } A;$
 $T \in \text{States } SA1 \rrbracket \implies$
 $\text{ran } (\text{CompFun } ((A \text{ [++]} (S,SA1) \text{ [++]} (T,SA2))) =$
 $\text{insert } \{\} (\text{insert } \{SA2\} (\text{ran } (\text{CompFun } (A \text{ [++]} (S,SA1))))))$
apply (*simp add: AddSA-HAStates AddSA-CompFun*)
apply (*subst FAddSA-ran*)
apply (*rule ballI*)
apply (*rule impI*)
apply (*subst AddSA-CompFun [THEN sym]*)
apply *simp*
apply *simp*
apply (*subst AddSA-CompFun [THEN sym]*)
apply *simp*
apply *simp*
apply (*rule CompFun-Int-disjoint*)
apply *simp*
apply (*simp add: AddSA-HAStates*)
apply (*simp add: AddSA-HAStates*)
apply (*case-tac S \in States SA1*)
apply *simp*
apply (*simp only: dom-CompFun [THEN sym]*)
apply (*frule FAddSA-dom-dom-States*)
apply *fast*

```

apply simp
apply (case-tac  $S \in \text{States } SA1$ )
apply simp
apply fast
apply (subst FAddSA-dom-dom-States)
apply simp
apply simp
apply simp
apply (case-tac  $S \in \text{States } SA1$ )
apply simp
apply fast
apply (subst FAddSA-dom-dom-States)
apply simp
apply simp
apply simp
apply (case-tac  $S \in \text{States } SA1$ )
apply simp
apply fast
apply simp
apply fast
done

```

lemma *AddSA-CompFun-ran-not-mem:*

```

 $\llbracket \text{States } SA2 \cap (\text{HStates } A \cup \text{States } SA1) = \{\};$ 
 $\text{States } SA1 \cap \text{HStates } A = \{\};$ 
 $S \in \text{HStates } A \rrbracket \implies$ 
 $\{SA2\} \notin \text{ran } (\text{CompFun } A [f+] (S, SA1))$ 
apply (cut-tac  $HA=A [++] (S, SA1)$  and  $Sas=\{SA2\}$  in ran-CompFun-is-not-SA)
apply (metis AddSA-HStates SA-States-disjunct2 SAs-States-HStates insert-subset)
apply (simp add: AddSA-HStates AddSA-CompFun)
done

```

lemma *AddSA-CompFun-ran3:*

```

 $\llbracket (\text{States } SA1 \cap \text{HStates } A) = \{\};$ 
 $(\text{States } SA2 \cap (\text{HStates } A \cup \text{States } SA1)) = \{\};$ 
 $(\text{States } SA3 \cap (\text{HStates } A \cup \text{States } SA1 \cup \text{States } SA2)) = \{\};$ 
 $S \in \text{HStates } A;$ 
 $T \in \text{States } SA1 \rrbracket \implies$ 
 $\text{ran } (\text{CompFun } ((A [++] (S, SA1)) [++] (T, SA2) [++] (T, SA3))) =$ 
 $\text{insert } \{\} (\text{insert } \{SA3, SA2\} (\text{ran } (\text{CompFun } (A [++] (S, SA1))))))$ 
apply (simp add: AddSA-HStates AddSA-CompFun)
apply (subst FAddSA-ran)
apply (metis AddSA-CompFun AddSA-HStates CompFun-Int-disjoint UnCI
dom-CompFun)
apply (metis AddSA-CompFun AddSA-HStates UnCI dom-CompFun)
apply (metis AddSA-CompFun AddSA-HStates UnCI dom-CompFun)

apply (subst AddSA-CompFun [THEN sym])
back

```

```

    apply simp
  apply simp

  apply (subst AddSA-CompFun [THEN sym])
    back
    apply (simp add: AddSA-HAStates)
    apply (simp add: AddSA-HAStates)
  apply (subst AddSA-CompFun-ran2)
    apply fast
    apply fast
    apply fast
    apply fast
  apply (simp add: AddSA-CompFun)
  apply (subst FAddSA-dom-insert)
    apply (subst FAddSA-dom-dom-States)
      apply simp
      apply fast
      apply simp
      apply fast
    apply (subst FAddSA-dom-emptyset)
      apply simp
      apply fast
      apply simp
      apply simp
    apply (subst FAddSA-dom-insert)
      apply (subst FAddSA-dom-dom-States)
        apply simp
        apply fast
        apply simp
        apply fast
      apply (subst FAddSA-dom-emptyset)
        apply simp
        apply fast
        apply simp
        apply simp
    by (simp add: AddSA-CompFun-ran-not-mem insert-Diff-if insert-commute)

```

lemma *AddSA-CompFun-PseudoHA-ran:*

```

  [ S ∈ States RootSA;
    States RootSA ∩ States SA = {} ] ⇒
  (ran (CompFun ((PseudoHA RootSA D) [++] (S,SA)))) = (insert {} {{SA}})

```

apply (subst AddSA-CompFun-ran)

apply *auto*

done

lemma *AddSA-CompFun-PseudoHA-ran2:*

```

  [ States SA1 ∩ States RootSA = {};
    States SA2 ∩ (States RootSA ∪ States SA1) = {};
    S ∈ States RootSA ] ⇒

```

```

      (ran (CompFun ((PseudoHA RootSA D) [++] (S,SA1) [++] (S,SA2)))) =
(insert {} {{SA2,SA1}})
apply (subst AddSA-CompFun-ran)
prefer 3
apply (subst AddSA-CompFun-the)
apply simp
apply simp
apply (subst AddSA-CompFun-PseudoHA-ran)
apply fast
apply fast
apply (subst AddSA-CompFun-the)
apply simp
apply simp
apply simp
apply fast
apply (simp add: AddSA-HAStates)
apply (simp add: AddSA-HAStates)
done

```

```

lemma AddSA-HAInitStates [simp]:
  [[ States SA  $\cap$  HAStates A = {}];
   S  $\in$  HAStates A ]  $\implies$ 
  HAINitStates (A [++] (S,SA)) = insert (InitState SA) (HAINitStates A)
apply (unfold HAINitStates-def)
apply (simp add: AddSA-SAs)
done

```

```

lemma AddSA-HAInitState [simp]:
  [[ States SA  $\cap$  HAStates A = {}];
   S  $\in$  HAStates A ]  $\implies$ 
  HAINitState (A [++] (S,SA)) = (HAINitState A)
apply (unfold HAINitState-def)
apply (simp add: AddSA-HARoot)
done

```

```

lemma AddSA-Chi [simp]:
  [[ States SA  $\cap$  HAStates A = {}];
   S  $\in$  HAStates A ]  $\implies$ 
  Chi (A [++] (S,SA)) S = (States SA)  $\cup$  (Chi A S)
apply (unfold Chi-def restrict-def)
apply (simp add: AddSA-SAs AddSA-HAStates AddSA-CompFun-the)
apply auto
done

```

```

lemma AddSA-Chi2 [simp]:
  [[ States SA  $\cap$  HAStates A = {}];
   S  $\in$  HAStates A;
   T  $\in$  States SA ]  $\implies$ 
  Chi (A [++] (S,SA)) T = {}

```


apply (*unfold Chi-def restrict-def*)
apply (*simp add: AddSA-SAs AddSA-HAStates AddSA-CompFun-the2*)
done

lemma *AddSA-Chi3* [*simp*]:
 $\llbracket \text{States } SA \cap \text{HAStates } A = \{\};$
 $S \in \text{HAStates } A;$
 $T \notin \text{States } SA; T \neq S \rrbracket \implies$
 $\text{Chi } (A \text{ [++]} (S, SA)) T = \text{Chi } A T$
apply (*unfold Chi-def restrict-def*)
apply (*simp add: AddSA-SAs AddSA-HAStates AddSA-CompFun-the3*)
apply *auto*
done

lemma *AddSA-ChiRel* [*simp*]:
 $\llbracket \text{States } SA \cap \text{HAStates } A = \{\};$
 $S \in \text{HAStates } A \rrbracket \implies$
 $\text{ChiRel } (A \text{ [++]} (S, SA)) = \{ (T, T') . T = S \wedge T' \in \text{States } SA \} \cup (\text{ChiRel } A)$
apply (*unfold ChiRel-def*)
apply (*simp add: AddSA-HAStates*)
apply *safe*
apply (*rename-tac T U*)
apply (*case-tac T \in States SA*)
apply *simp*
apply *simp*
apply (*rename-tac T U*)
apply (*case-tac T \neq S*)
apply (*case-tac T \in States SA*)
apply *simp*
apply *simp*
apply *simp*
apply (*rename-tac T U*)
apply (*case-tac T \in States SA*)
apply *simp*
apply *simp*
apply (*cut-tac A=A and T=T in Chi-HAStates*)
apply *fast*
apply (*case-tac T \in States SA*)
apply *simp*
apply *simp*
apply (*cut-tac A=A and T=T in Chi-HAStates*)
apply *fast*
apply *fast*
apply (*rename-tac T U*)
apply (*case-tac T \neq S*)
apply (*case-tac T \in States SA*)
apply *simp*
apply *simp*
apply *simp*

apply (*rename-tac* $T\ U$)
apply (*case-tac* $T \in \text{States } SA$)
apply *auto*
apply (*metis* *AddSA-Chi* *AddSA-Chi3* *Int-iff* *Un-iff* *empty-iff*)
done

lemma *help-InitConf*:

$\llbracket \text{States } SA \cap \text{HAStates } A = \{\} \rrbracket \implies \{p. \text{fst } p \neq \text{InitState } SA \wedge \text{snd } p \neq \text{InitState } SA \wedge$

$p \in \text{insert } (\text{InitState } SA) (\text{HAIInitStates } A) \times \text{insert } (\text{InitState } SA) (\text{HAIInitStates } A) \wedge$

$(p \in \{S\} \times \text{States } SA \vee p \in \text{ChiRel } A)\} =$

$(\text{HAIInitStates } A \times \text{HAIInitStates } A \cap \text{ChiRel } A)$

apply *auto*

apply (*cut-tac* $A=SA$ **in** *InitState-States*)

apply (*cut-tac* $A=A$ **in** *HAIInitStates-HAStates*, *fast*)

apply (*cut-tac* $A=SA$ **in** *InitState-States*)

apply (*cut-tac* $A=A$ **in** *HAIInitStates-HAStates*, *fast*)

done

lemma *AddSA-InitConf* [*simp*]:

$\llbracket \text{States } SA \cap \text{HAStates } A = \{\};$

$S \in \text{InitConf } A \rrbracket \implies$

$\text{InitConf } (A \text{ [++]} (S, SA)) = \text{insert } (\text{InitState } SA) (\text{InitConf } A)$

apply (*frule* *InitConf-HAStates2*)

apply (*unfold* *InitConf-def*)

apply (*simp del*: *insert-Times-insert*)

apply *auto*

apply (*rename-tac* T)

apply (*case-tac* $T=S$)

apply *auto*

prefer 3

apply (*rule-tac* $R=(\text{HAIInitStates } A) \times (\text{HAIInitStates } A) \cap \text{ChiRel } A$ **in** *trancl-subseteq*)

apply *auto*

apply (*rotate-tac* 3)

apply (*frule* *trancl-collect*)

prefer 2

apply *fast*

apply *auto*

apply (*cut-tac* $A=SA$ **in** *InitState-States*)

apply (*frule* *ChiRel-HAStates*)

apply *fast*

apply (*frule* *ChiRel-HAStates*)

apply (*cut-tac* $A=SA$ **in** *InitState-States*)

apply *fast*

apply (*frule* *ChiRel-HAStates*)

apply (*cut-tac* $A=SA$ **in** *InitState-States*)

apply *fast*

apply (*subst* *help-InitConf* [*THEN sym*])

```

apply fast
apply auto
apply (rule-tac  $b=S$  in rtrancl-into-rtrancl)
apply auto
prefer 2
apply (erule rtranclE)
apply auto
prefer 2
apply (erule rtranclE)
apply auto
apply (rule-tac  $R=(H\text{InitStates } A) \times (H\text{InitStates } A) \cap \text{ChiRel } A$  in trancl-subseteq)
apply auto
done

```

```

lemma AddSA-InitConf2 [simp]:
   $\llbracket \text{States } SA \cap H\text{States } A = \{\};$ 
   $S \notin \text{InitConf } A;$ 
   $S \in H\text{States } A \rrbracket \implies$ 
   $\text{InitConf } (A \text{ [++]} (S,SA)) = \text{InitConf } A$ 
apply (unfold InitConf-def)
apply simp
apply auto
apply (rename-tac T)
prefer 2
apply (rule-tac  $R=(H\text{InitStates } A) \times (H\text{InitStates } A) \cap \text{ChiRel } A$  in trancl-subseteq)
apply auto
apply (case-tac  $T=\text{InitState } SA$ )
apply auto
prefer 2
apply (rotate-tac 3)
apply (frule trancl-collect)
prefer 2
apply fast
apply auto
apply (cut-tac  $A=SA$  in InitState-States)
apply (frule ChiRel-HAStates)
apply fast
apply (cut-tac  $A=SA$  in InitState-States)
apply (frule ChiRel-HAStates)
apply fast
apply (cut-tac  $A=SA$  in InitState-States)
apply (cut-tac  $A=A$  in H\text{InitStates-HAStates})
apply fast
apply (subst help-InitConf [THEN sym])
apply fast
apply auto
apply (rule-tac  $b=\text{InitState } SA$  in rtrancl-induct)
apply auto
apply (frule ChiRel-HAStates2)

```

```

apply (cut-tac  $A=SA$  in InitState-States)
apply fast
prefer 2
apply (frule ChiRel-HAStates)
apply (cut-tac  $A=SA$  in InitState-States)
apply fast
apply (rule rtrancl-into-rtrancl)
apply auto
apply (rule rtrancl-into-rtrancl)
apply auto
done

```

11.5 Theorems for Calculating Wellformedness of HA

lemma *PseudoHA-HAStates-IFF*:

```

(States  $SA$ ) =  $X$   $\implies$  (HAStates (PseudoHA  $SA$   $D$ )) =  $X$ 
apply simp
done

```

lemma *AddSA-SAs-IFF*:

```

 $\llbracket$  States  $SA \cap$  HAStates  $HA = \{\}$ ;
 $S \in$  HAStates  $HA$ ;
(SAs  $HA$ ) =  $X$   $\rrbracket \implies$  (SAs ( $HA$   $[[+]]$  ( $S, SA$ ))) = (insert  $SA$   $X$ )
apply (subst AddSA-SAs)
apply auto
done

```

lemma *AddSA-Events-IFF*:

```

 $\llbracket$  States  $SA \cap$  HAStates  $HA = \{\}$ ;
 $S \in$  HAStates  $HA$ ;
(HAEvents  $HA$ ) =  $HAE$ ;
(SAEvents  $SA$ ) =  $SAE$ ;
( $HAE \cup$   $SAE$ ) =  $X$   $\rrbracket \implies$  (HAEvents ( $HA$   $[[+]]$  ( $S, SA$ ))) =  $X$ 
apply (subst AddSA-Events)
apply auto
done

```

lemma *AddSA-CompFun-IFF*:

```

 $\llbracket$  States  $SA \cap$  HAStates  $HA = \{\}$ ;
 $S \in$  HAStates  $HA$ ;
(CompFun  $HA$ ) =  $HAG$ ;
( $HAG$   $[f+]$  ( $S, SA$ )) =  $X$   $\rrbracket \implies$  (CompFun ( $HA$   $[[+]]$  ( $S, SA$ ))) =  $X$ 
apply (subst AddSA-CompFun)
apply auto
done

```

lemma *AddSA-HAStates-IFF*:

```

 $\llbracket$  States  $SA \cap$  HAStates  $HA = \{\}$ ;
 $S \in$  HAStates  $HA$ ;

```

$(HAStates HA) = HAS;$
 $(States SA) = SAS;$
 $(HAS \cup SAS) = X \implies (HAStates (HA [++] (S, SA))) = X$
apply (subst AddSA-HAStates)
apply auto
done

lemma AddSA-HAInitValue-IFF:
 $\llbracket States SA \cap HAStates HA = \{\};$
 $S \in HAStates HA;$
 $(HAInitValue HA) = X \implies (HAInitValue (HA [++] (S, SA))) = X$
apply (subst AddSA-HAInitValue)
apply auto
done

lemma AddSA-HARoot-IFF:
 $\llbracket States SA \cap HAStates HA = \{\};$
 $S \in HAStates HA;$
 $(HARoot HA) = X \implies (HARoot (HA [++] (S, SA))) = X$
apply (subst AddSA-HARoot)
apply auto
done

lemma AddSA-InitConf-IFF:
 $\llbracket InitConf A = Y;$
 $States SA \cap HAStates A = \{\};$
 $S \in HAStates A;$
 $(if S \in Y then insert (InitState SA) Y else Y) = X \implies$
 $InitConf (A [++] (S, SA)) = X$
apply (case-tac $S \in Y$)
apply auto
done

lemma AddSA-CompFun-ran-IFF:
 $\llbracket (States SA \cap HAStates A) = \{\};$
 $S \in HAStates A;$
 $(insert \{\} (insert (insert SA (the ((CompFun A) S))) (ran (CompFun A) -$
 $\{the ((CompFun A) S)\})) = X \implies$
 $ran (CompFun (A [++] (S, SA))) = X$
apply (subst AddSA-CompFun-ran)
apply auto
done

lemma AddSA-CompFun-ran2-IFF:
 $\llbracket (States SA1 \cap HAStates A) = \{\};$
 $(States SA2 \cap (HAStates A \cup States SA1)) = \{\};$
 $S \in HAStates A;$
 $T \in States SA1;$
 $insert \{\} (insert \{SA2\} (ran (CompFun (A [++] (S, SA1)))))) = X \implies$

```

    ran (CompFun ((A [++] (S,SA1)) [++] (T,SA2))) = X
apply (subst AddSA-CompFun-ran2)
apply auto
done

```

lemma AddSA-CompFun-ran3-IFF:

```

[[ (States SA1 ∩ HAStates A) = {};
   (States SA2 ∩ (HAStates A ∪ States SA1)) = {};
   (States SA3 ∩ (HAStates A ∪ States SA1 ∪ States SA2)) = {};
   S ∈ HAStates A;
   T ∈ States SA1;
   insert {} (insert {SA3,SA2} (ran (CompFun (A [++] (S,SA1)))))) = X ]] ⇒
   ran (CompFun ((A [++] (S,SA1)) [++] (T,SA2) [++] (T,SA3))) = X
apply (subst AddSA-CompFun-ran3)
apply auto
done

```

lemma AddSA-CompFun-PseudoHA-ran-IFF:

```

[[ S ∈ States RootSA;
   States RootSA ∩ States SA = {};
   (insert {} {{SA}}) = X ]] ⇒
   ran (CompFun ((PseudoHA RootSA D) [++] (S,SA))) = X
apply (subst AddSA-CompFun-PseudoHA-ran)
apply auto
done

```

lemma AddSA-CompFun-PseudoHA-ran2-IFF:

```

[[ States SA1 ∩ States RootSA = {};
   States SA2 ∩ (States RootSA ∪ States SA1) = {};
   S ∈ States RootSA;
   (insert {} {{SA2,SA1}}) = X ]] ⇒
   ran (CompFun ((PseudoHA RootSA D) [++] (S,SA1) [++] (S,SA2))) = X
apply (subst AddSA-CompFun-PseudoHA-ran2)
apply auto
done

```

ML ‹

```

val AddSA-SAs-IFF = @{thm AddSA-SAs-IFF};
val AddSA-Events-IFF = @{thm AddSA-Events-IFF};
val AddSA-CompFun-IFF = @{thm AddSA-CompFun-IFF};
val AddSA-HAStates-IFF = @{thm AddSA-HAStates-IFF};
val PseudoHA-HAStates-IFF = @{thm PseudoHA-HAStates-IFF};
val AddSA-HAInitValue-IFF = @{thm AddSA-HAInitValue-IFF};
val AddSA-CompFun-ran-IFF = @{thm AddSA-CompFun-ran-IFF};
val AddSA-HARoot-IFF = @{thm AddSA-HARoot-IFF};
val insert-inter = @{thm insert-inter};
val insert-notmem = @{thm insert-notmem};

```

```

val PseudoHA-CompFun = @{thm PseudoHA-CompFun};
val PseudoHA-Events = @{thm PseudoHA-Events};
val PseudoHA-SAs = @{thm PseudoHA-SAs};
val PseudoHA-HARoot = @{thm PseudoHA-HARoot};
val PseudoHA-HAInitValue = @{thm PseudoHA-HAInitValue};
val PseudoHA-CompFun-ran = @{thm PseudoHA-CompFun-ran};
val Un-empty-right = @{thm Un-empty-right};
val insert-union = @{thm insert-union};

```

```

fun wellformed-tac ctxt L i =
  FIRST[resolve-tac ctxt [AddSA-SAs-IFF] i,
        resolve-tac ctxt [AddSA-Events-IFF] i,
        resolve-tac ctxt [AddSA-CompFun-IFF] i,
        resolve-tac ctxt [AddSA-HAStates-IFF] i,
        resolve-tac ctxt [PseudoHA-HAStates-IFF] i,
        resolve-tac ctxt [AddSA-HAInitValue-IFF] i,
        resolve-tac ctxt [AddSA-HARoot-IFF] i,
        resolve-tac ctxt [AddSA-CompFun-ran-IFF] i,
        resolve-tac ctxt [insert-inter] i,
        resolve-tac ctxt [insert-notmem] i,
        CHANGED (simp-tac (put-simpset HOL-basic-ss ctxt addsimps
[PseudoHA-HARoot, PseudoHA-CompFun, PseudoHA-CompFun-ran, PseudoHA-Events, PseudoHA-SAs,
PseudoHA-HAInitValue, Un-empty-right]@ L) i),
        fast-tac ctxt i,
        CHANGED (simp-tac ctxt i)];
)

```

```

method-setup wellformed = ⟨Attrib.thms⟩⟩ (fn thms => fn ctxt => (METHOD
(fn facts =>
  (HEADGOAL (wellformed-tac ctxt (facts @ thms))))))

```

end

12 Example Specification for a Car Audio System

```

theory CarAudioSystem
imports HAKripke HAOps
begin

```

12.1 Definitions

12.1.1 Data space for two Integer-Variables

```

datatype d = V0 int
          | V1 int

```

```

primrec
  Sel0 :: d => int where

```

$Sel0 (V0 i) = i$

primrec

$Sel1 :: d \Rightarrow int$ **where**
 $Sel1 (V1 i) = i$

definition

$Select0 :: [d data] \Rightarrow int$ **where**
 $Select0 d = Sel0 (DataPart d 0)$

definition

$Select1 :: [d data] \Rightarrow int$ **where**
 $Select1 d = Sel1 (DataPart d 1)$

definition

$DSpace :: d dataspace$ **where**
 $DSpace = Abs-dataspace ([range V0, range V1])$

definition

$LiftInitData :: (d list) \Rightarrow d data$ **where**
 $LiftInitData L = Abs-data (L, DSpace)$

definition

$LiftPUpdate :: (d data \Rightarrow ((d option) list)) \Rightarrow d pupdate$ **where**
 $LiftPUpdate L = Abs-pupdate$
 $(\lambda d. \text{if } ((DataSpace d) = DSpace) \text{ then}$
 $\quad Abs-pdata (L d, DSpace)$
 $\quad \text{else } (Data2PData d))$

12.1.2 Sequential Automaton *Root-CTRL*

definition

$Root-CTRL-States :: string set$ **where**
 $Root-CTRL-States = \{ "CarAudioSystem" \}$

definition

$Root-CTRL-Init :: string$ **where**
 $Root-CTRL-Init = "CarAudioSystem"$

definition

$Root-CTRL-Labels :: (string, string, d) label set$ **where**
 $Root-CTRL-Labels = \{ \}$

definition

Root-CTRL-Delta :: (string,string,d)trans set **where**
Root-CTRL-Delta = {}

definition

Root-CTRL :: (string,string,d)seqauto **where**
Root-CTRL = Abs-seqauto (*Root-CTRL-States*, *Root-CTRL-Init*,
Root-CTRL-Labels, *Root-CTRL-Delta*)

12.1.3 Sequential Automaton CDPlayer-CTRL

definition

CDPlayer-CTRL-States :: string set **where**
CDPlayer-CTRL-States = {"ReadTracks","CDFull","CDEmpty"}

definition

CDPlayer-CTRL-Init :: string **where**
CDPlayer-CTRL-Init = "CDEmpty"

definition

CDPlayer-CTRL-Update1 :: d pupdate **where**
CDPlayer-CTRL-Update1 = LiftPUpdate (% d. [None, Some (V1 ((Select0 d) + 1))])

definition

CDPlayer-CTRL-Action1 :: (string,d)action **where**
CDPlayer-CTRL-Action1 = ({} , *CDPlayer-CTRL-Update1*)

definition

CDPlayer-CTRL-Update2 :: d pupdate **where**
CDPlayer-CTRL-Update2 = LiftPUpdate (% d. [Some (V0 ((Select0 d) + 1)), None])

definition

CDPlayer-CTRL-Action2 :: (string,d)action **where**
CDPlayer-CTRL-Action2 = ({} , *CDPlayer-CTRL-Update2*)

definition

CDPlayer-CTRL-Update3 :: d pupdate **where**
CDPlayer-CTRL-Update3 = LiftPUpdate (% d. [Some (V0 0), None])

definition

CDPlayer-CTRL-Action3 :: (string,d)action **where**
CDPlayer-CTRL-Action3 = ({} , *CDPlayer-CTRL-Update3*)

definition

CDPlayer-CTRL-Labels :: (string,string,d)label set **where**
CDPlayer-CTRL-Labels = {(En "LastTrack", defaultguard, *CDPlayer-CTRL-Action1*),
(En "NewTrack", defaultguard, *CDPlayer-CTRL-Action2*),
(And (En "CDEject") (In "On"), defaultguard, CD-

Player-CTRL-Action3),
 (En "CDIn", defaultguard, defaultaction)}

definition

CDPlayer-CTRL-Delta :: (string,string,d)trans set **where**
CDPlayer-CTRL-Delta = {"ReadTracks", (En "LastTrack", defaultguard, CD-
Player-CTRL-Action1),
 "CDFull"},
 ("CDFull", (And (En "CDEject") (In "On"), defaultguard,
CDPlayer-CTRL-Action3),
 "CDEmpty"),
 ("ReadTracks", (En "NewTrack", defaultguard, CD-
Player-CTRL-Action2),
 "ReadTracks"),
 ("CDEmpty", (En "CDIn", defaultguard, defaultaction),
 "ReadTracks")}

definition

CDPlayer-CTRL :: (string,string,d)seqauto **where**
CDPlayer-CTRL = Abs-seqauto (*CDPlayer-CTRL-States*, *CDPlayer-CTRL-Init*,
CDPlayer-CTRL-Labels, *CDPlayer-CTRL-Delta*)

12.1.4 Sequential Automaton *AudioPlayer-CTRL*

definition

AudioPlayer-CTRL-States :: string set **where**
AudioPlayer-CTRL-States = {"Off", "On"}

definition

AudioPlayer-CTRL-Init :: string **where**
AudioPlayer-CTRL-Init = "Off"

definition

AudioPlayer-CTRL-Labels :: (string,string,d)label set **where**
AudioPlayer-CTRL-Labels = {(En "O", defaultguard, defaultaction)}

definition

AudioPlayer-CTRL-Delta :: (string,string,d)trans set **where**
AudioPlayer-CTRL-Delta = {"Off", (En "O", defaultguard, defaultaction), "On"},
 ("On", (En "O", defaultguard, defaultaction), "Off")}

definition

AudioPlayer-CTRL :: (string,string,d)seqauto **where**
AudioPlayer-CTRL = Abs-seqauto (*AudioPlayer-CTRL-States*, *AudioPlayer-CTRL-Init*,
AudioPlayer-CTRL-Labels, *AudioPlayer-CTRL-Delta*)

12.1.5 Sequential Automaton *On-CTRL*

definition

On-CTRL-States :: string set **where**

$On-CTRL-States = \{ "TunerMode", "CDMode" \}$

definition

$On-CTRL-Init :: string$ **where**
 $On-CTRL-Init = "TunerMode"$

definition

$On-CTRL-Labels :: (string, string, d)label$ set **where**
 $On-CTRL-Labels = \{ (And (En "Src") (In "CDFull")), defaultguard, defaultaction),$
 $(En "Src", defaultguard, defaultaction),$
 $(En "CDEject", defaultguard, defaultaction),$
 $(En "EndOfTitle", (\lambda d. (DataPart d 0) = (DataPart d 1)),$
 $defaultaction) \}$

definition

$On-CTRL-Delta :: (string, string, d)trans$ set **where**
 $On-CTRL-Delta = \{ ("TunerMode", (And (En "Src") (In "CDFull")), defaultguard, defaultaction), "CDMode"),$
 $("CDMode", (En "Src", defaultguard, defaultaction), "TunerMode"),$
 $("CDMode", (En "CDEject", defaultguard, defaultaction),$
 $"TunerMode"),$
 $("CDMode", (En "EndOfTitle", (\lambda d. (DataPart d 0) = (DataPart d 1)), defaultaction),$
 $"TunerMode") \}$

definition

$On-CTRL :: (string, string, d)seqauto$ **where**
 $On-CTRL = Abs-seqauto (On-CTRL-States, On-CTRL-Init,$
 $On-CTRL-Labels, On-CTRL-Delta)$

12.1.6 Sequential Automaton *TunerMode-CTRL*

definition

$TunerMode-CTRL-States :: string$ set **where**
 $TunerMode-CTRL-States = \{ "1", "2", "3", "4" \}$

definition

$TunerMode-CTRL-Init :: string$ **where**
 $TunerMode-CTRL-Init = "1"$

definition

$TunerMode-CTRL-Labels :: (string, string, d)label$ set **where**
 $TunerMode-CTRL-Labels = \{ (En "Next", defaultguard, defaultaction),$
 $(En "Back", defaultguard, defaultaction) \}$

definition

$TunerMode-CTRL-Delta :: (string, string, d)trans$ set **where**
 $TunerMode-CTRL-Delta = \{ ("1", (En "Next", defaultguard, defaultaction), "2"),$

```

("2", (En "Next", defaultguard, defaultaction), "3"),
("3", (En "Next", defaultguard, defaultaction), "4"),
("4", (En "Next", defaultguard, defaultaction), "1"),
("1", (En "Back", defaultguard, defaultaction), "4"),
("4", (En "Back", defaultguard, defaultaction), "3"),
("3", (En "Back", defaultguard, defaultaction), "2"),
("2", (En "Back", defaultguard, defaultaction), "1")}

```

definition

```

TunerMode-CTRL :: (string,string,d)seqauto where
TunerMode-CTRL = Abs-seqauto (TunerMode-CTRL-States, TunerMode-CTRL-Init,
                             TunerMode-CTRL-Labels, TunerMode-CTRL-Delta)

```

12.1.7 Sequential Automaton *CDMode-CTRL*

definition

```

CDMode-CTRL-States :: string set where
CDMode-CTRL-States = {"Playing", "SelectingNextTrack",
                      "SelectingPreviousTrack"}

```

definition

```

CDMode-CTRL-Init :: string where
CDMode-CTRL-Init = "Playing"

```

definition

```

CDMode-CTRL-Update1 :: d pupdate where
CDMode-CTRL-Update1 = LiftPUpdate (% d. [ Some (V0 ((Select0 d) + 1)),
None ])

```

definition

```

CDMode-CTRL-Action1 :: (string,d)action where
CDMode-CTRL-Action1 = ({}, CDMode-CTRL-Update1)

```

definition

```

CDMode-CTRL-Update2 :: d pupdate where
CDMode-CTRL-Update2 = LiftPUpdate (% d. [ Some (V0 ((Select0 d) - 1)),
None ])

```

definition

```

CDMode-CTRL-Action2 :: (string,d)action where
CDMode-CTRL-Action2 = ({}, CDMode-CTRL-Update2)

```

definition

```

CDMode-CTRL-Labels :: (string,string,d)label set where
CDMode-CTRL-Labels = {(En "Next", defaultguard, defaultaction),
                      (En "Back", defaultguard, defaultaction),
                      (En "Ready", defaultguard, CDMode-CTRL-Action1),
                      (En "Ready", defaultguard, CDMode-CTRL-Action2),
                      (En "EndOfTitle", (\ (d:: d data). (Select0 d) < (Select1 d)),

```

defaultaction)}

definition

CDMode-CTRL-Delta :: (*string, string, d*)*trans set* **where**
CDMode-CTRL-Delta = {"Playing", (En "Next", defaultguard, defaultaction) ,
 "SelectingNextTrack"),
 ("SelectingNextTrack", (En "Ready", defaultguard, CD-
 Mode-CTRL-Action1) ,"Playing"),
 ("Playing", (En "Back", defaultguard, defaultaction)
 ,"SelectingPreviousTrack"),
 ("SelectingPreviousTrack", (En "Ready", defaultguard,
 CDMode-CTRL-Action2),
 "Playing"),
 ("Playing", (En "EndOfTitle", (λ (*d*:: *d data*). (Select0 *d*) <
 (Select1 *d*)), defaultaction),
 "SelectingNextTrack")}

definition

CDMode-CTRL :: (*string, string, d*)*seqauto* **where**
CDMode-CTRL = *Abs-seqauto* (*CDMode-CTRL-States*, *CDMode-CTRL-Init*,
CDMode-CTRL-Labels, *CDMode-CTRL-Delta*)

12.1.8 Hierarchical Automaton *CarAudioSystem*

definition

CarAudioSystem :: (*string, string, d*)*hierauto* **where**
CarAudioSystem = ((*PseudoHA Root-CTRL* (*LiftInitData* [*V0 0*, *V1 0*]))
 [++] ("*CarAudioSystem*", *CDPlayer-CTRL*)
 [++] ("*CarAudioSystem*", *AudioPlayer-CTRL*)
 [++] ("*On*", *TunerMode-CTRL*)
 [++] ("*On*", *CDMode-CTRL*))

12.2 Lemmas

12.2.1 Sequential Automaton *CDMode-CTRL*

lemma *check-Root-CTRL*:

(*Root-CTRL-States*, *Root-CTRL-Init*, *Root-CTRL-Labels*, *Root-CTRL-Delta*) : *seqauto*

apply (*unfold seqauto-def SeqAuto-def Root-CTRL-States-def Root-CTRL-Init-def Root-CTRL-Labels-def Root-CTRL-Delta-def*)

apply *simp*

done

lemma *States-Root-CTRL*:

States Root-CTRL = *Root-CTRL-States*

apply (*simp add: Root-CTRL-def*)

apply (*unfold States-def*)

apply (*simp add: Abs-seqauto-inverse check-Root-CTRL*)

done

lemma *Init-State-Root-CTRL*:
InitState Root-CTRL = Root-CTRL-Init
apply (*simp add: Root-CTRL-def*)
apply (*unfold InitState-def*)
apply (*simp add: Abs-seqauto-inverse check-Root-CTRL*)
done

lemma *Labels-Root-CTRL*:
Labels Root-CTRL = Root-CTRL-Labels
apply (*simp add: Root-CTRL-def*)
apply (*unfold Labels-def*)
apply (*simp add: Abs-seqauto-inverse check-Root-CTRL*)
done

lemma *Delta-Root-CTRL*:
Delta Root-CTRL = Root-CTRL-Delta
apply (*simp add: Root-CTRL-def*)
apply (*unfold Delta-def*)
apply (*simp add: Abs-seqauto-inverse check-Root-CTRL*)
done

schematic-goal *Events-Root-CTRL*:
SAEvents Root-CTRL = ?X
apply (*unfold SAEvents-def expr-def*)
apply (*rule trans*)
apply (*simp add: expr-def Delta-Root-CTRL Root-CTRL-Delta-def*)
apply (*rule refl*)
done

12.2.2 Sequential Automaton *CDPlayer-CTRL*

lemma *check-CDPlayer-CTRL*:
(*CDPlayer-CTRL-States, CDPlayer-CTRL-Init, CDPlayer-CTRL-Labels, CDPlayer-CTRL-Delta*)
: *seqauto*
apply (*unfold seqauto-def SeqAuto-def CDPlayer-CTRL-States-def CDPlayer-CTRL-Init-def*
CDPlayer-CTRL-Labels-def CDPlayer-CTRL-Delta-def)
apply *simp*
done

lemma *States-CDPlayer-CTRL*:
States CDPlayer-CTRL = CDPlayer-CTRL-States
apply (*simp add: CDPlayer-CTRL-def*)
apply (*unfold States-def*)
apply (*simp add: Abs-seqauto-inverse check-CDPlayer-CTRL*)
done

lemma *Init-State-CDPlayer-CTRL*:
InitState CDPlayer-CTRL = CDPlayer-CTRL-Init

apply (*simp add: CDPlayer-CTRL-def*)
apply (*unfold initState-def*)
apply (*simp add: Abs-seqauto-inverse check-CDPlayer-CTRL*)
done

lemma *Labels-CDPlayer-CTRL:*
Labels CDPlayer-CTRL = CDPlayer-CTRL-Labels
apply (*simp add: CDPlayer-CTRL-def*)
apply (*unfold Labels-def*)
apply (*simp add: Abs-seqauto-inverse check-CDPlayer-CTRL*)
done

lemma *Delta-CDPlayer-CTRL:*
Delta CDPlayer-CTRL = CDPlayer-CTRL-Delta
apply (*simp add: CDPlayer-CTRL-def*)
apply (*unfold Delta-def*)
apply (*simp add: Abs-seqauto-inverse check-CDPlayer-CTRL*)
done

schematic-goal *Events-CDPlayer-CTRL:*
SAEvents CDPlayer-CTRL = ?X
apply (*unfold SAEvents-def*)
apply (*rule trans*)
apply (*simp add: expr-def Delta-CDPlayer-CTRL CDPlayer-CTRL-Delta-def CD-*
Player-CTRL-Action1-def CDPlayer-CTRL-Action2-def CDPlayer-CTRL-Action3-def
Label-def)
apply (*rule refl*)
done

12.2.3 Sequential Automaton *AudioPlayer-CTRL*

lemma *check-AudioPlayer-CTRL:*
(AudioPlayer-CTRL-States,AudioPlayer-CTRL-Init,AudioPlayer-CTRL-Labels,AudioPlayer-CTRL-Delta)
: seqauto
apply (*unfold seqauto-def SeqAuto-def AudioPlayer-CTRL-States-def AudioPlayer-CTRL-Init-def*
AudioPlayer-CTRL-Labels-def AudioPlayer-CTRL-Delta-def)
apply *simp*
done

lemma *States-AudioPlayer-CTRL:*
States AudioPlayer-CTRL = AudioPlayer-CTRL-States
apply (*simp add: AudioPlayer-CTRL-def*)
apply (*unfold States-def*)
apply (*simp add: Abs-seqauto-inverse check-AudioPlayer-CTRL*)
done

lemma *Init-State-AudioPlayer-CTRL:*
InitState AudioPlayer-CTRL = AudioPlayer-CTRL-Init
apply (*simp add: AudioPlayer-CTRL-def*)

apply (*unfold initState-def*)
apply (*simp add: Abs-seqauto-inverse check-AudioPlayer-CTRL*)
done

lemma *Labels-AudioPlayer-CTRL*:
Labels AudioPlayer-CTRL = AudioPlayer-CTRL-Labels
apply (*simp add: AudioPlayer-CTRL-def*)
apply (*unfold Labels-def*)
apply (*simp add: Abs-seqauto-inverse check-AudioPlayer-CTRL*)
done

lemma *Delta-AudioPlayer-CTRL*:
Delta AudioPlayer-CTRL = AudioPlayer-CTRL-Delta
apply (*simp add: AudioPlayer-CTRL-def*)
apply (*unfold Delta-def*)
apply (*simp add: Abs-seqauto-inverse check-AudioPlayer-CTRL*)
done

schematic-goal *Events-AudioPlayer-CTRL*:
SAEvents AudioPlayer-CTRL = ?X
apply (*unfold SAEvents-def*)
apply (*rule trans*)
apply (*simp add: expr-def Delta-AudioPlayer-CTRL AudioPlayer-CTRL-Delta-def Label-def*)
apply (*rule refl*)
done

12.2.4 Sequential Automaton *On-CTRL*

lemma *check-On-CTRL*:
(On-CTRL-States, On-CTRL-Init, On-CTRL-Labels, On-CTRL-Delta) : seqauto
apply (*unfold seqauto-def SeqAuto-def On-CTRL-States-def On-CTRL-Init-def On-CTRL-Labels-def On-CTRL-Delta-def*)
apply *simp*
done

lemma *States-On-CTRL*:
States On-CTRL = On-CTRL-States
apply (*simp add: On-CTRL-def*)
apply (*unfold States-def*)
apply (*simp add: Abs-seqauto-inverse check-On-CTRL*)
done

lemma *Init-State-On-CTRL*:
InitState On-CTRL = On-CTRL-Init
apply (*simp add: On-CTRL-def*)
apply (*unfold initState-def*)
apply (*simp add: Abs-seqauto-inverse check-On-CTRL*)
done

lemma *Labels-On-CTRL*:
 $Labels\ On-CTRL = On-CTRL-Labels$
apply (*simp add: On-CTRL-def*)
apply (*unfold Labels-def*)
apply (*simp add: Abs-seqauto-inverse check-On-CTRL*)
done

lemma *Delta-On-CTRL*:
 $Delta\ On-CTRL = On-CTRL-Delta$
apply (*simp add: On-CTRL-def*)
apply (*unfold Delta-def*)
apply (*simp add: Abs-seqauto-inverse check-On-CTRL*)
done

schematic-goal *Events-On-CTRL*:
 $SAEvents\ On-CTRL = ?X$
apply (*unfold SAEvents-def*)
apply (*rule trans*)
apply (*simp add: expr-def Delta-On-CTRL On-CTRL-Delta-def Label-def*)
apply (*rule refl*)
done

12.2.5 Sequential Automaton *TunerMode-CTRL*

lemma *check-TunerMode-CTRL*:
(*TunerMode-CTRL-States, TunerMode-CTRL-Init, TunerMode-CTRL-Labels, TunerMode-CTRL-Delta*)
: *seqauto*
apply (*unfold seqauto-def SeqAuto-def TunerMode-CTRL-States-def TunerMode-CTRL-Init-def*
TunerMode-CTRL-Labels-def TunerMode-CTRL-Delta-def)
apply *simp*
done

lemma *States-TunerMode-CTRL*:
 $States\ TunerMode-CTRL = TunerMode-CTRL-States$
apply (*simp add: TunerMode-CTRL-def*)
apply (*unfold States-def*)
apply (*simp add: Abs-seqauto-inverse check-TunerMode-CTRL*)
done

lemma *Init-State-TunerMode-CTRL*:
 $InitState\ TunerMode-CTRL = TunerMode-CTRL-Init$
apply (*simp add: TunerMode-CTRL-def*)
apply (*unfold InitState-def*)
apply (*simp add: Abs-seqauto-inverse check-TunerMode-CTRL*)
done

lemma *Labels-TunerMode-CTRL*:
 $Labels\ TunerMode-CTRL = TunerMode-CTRL-Labels$

```

apply (simp add: TunerMode-CTRL-def)
apply (unfold Labels-def)
apply (simp add: Abs-seqauto-inverse check-TunerMode-CTRL)
done

```

```

lemma Delta-TunerMode-CTRL:
  Delta TunerMode-CTRL = TunerMode-CTRL-Delta
apply (simp add: TunerMode-CTRL-def)
apply (unfold Delta-def)
apply (simp add: Abs-seqauto-inverse check-TunerMode-CTRL)
done

```

```

schematic-goal Events-TunerMode-CTRL:
  SAEvents TunerMode-CTRL = ?X
apply (unfold SAEvents-def)
apply (rule trans)
apply (simp add: expr-def Delta-TunerMode-CTRL TunerMode-CTRL-Delta-def
  Label-def)
apply (rule refl)
done

```

12.2.6 Sequential Automaton *CDMode-CTRL*

```

lemma check-CDMode-CTRL:
  (CDMode-CTRL-States, CDMode-CTRL-Init, CDMode-CTRL-Labels, CDMode-CTRL-Delta)
  : seqauto
apply (unfold seqauto-def SeqAuto-def CDMode-CTRL-States-def CDMode-CTRL-Init-def
  CDMode-CTRL-Labels-def CDMode-CTRL-Delta-def)
apply simp
done

```

```

lemma States-CDMode-CTRL:
  States CDMode-CTRL = CDMode-CTRL-States
apply (simp add: CDMode-CTRL-def)
apply (unfold States-def)
apply (simp add: Abs-seqauto-inverse check-CDMode-CTRL)
done

```

```

lemma Init-State-CDMode-CTRL:
  InitState CDMode-CTRL = CDMode-CTRL-Init
apply (simp add: CDMode-CTRL-def)
apply (unfold InitState-def)
apply (simp add: Abs-seqauto-inverse check-CDMode-CTRL)
done

```

```

lemma Labels-CDMode-CTRL:
  Labels CDMode-CTRL = CDMode-CTRL-Labels
apply (simp add: CDMode-CTRL-def)
apply (unfold Labels-def)

```

apply (*simp add: Abs-seqauto-inverse check-CDMode-CTRL*)
done

lemma *Delta-CDMode-CTRL*:
Delta CDMode-CTRL = CDMode-CTRL-Delta
apply (*simp add: CDMode-CTRL-def*)
apply (*unfold Delta-def*)
apply (*simp add: Abs-seqauto-inverse check-CDMode-CTRL*)
done

schematic-goal *Events-CDMode-CTRL*:
SAEvents CDMode-CTRL = ?X
apply (*unfold SAEvents-def*)
apply (*rule trans*)
apply (*simp add: expr-def Label-def Delta-CDMode-CTRL CDMode-CTRL-Delta-def*
CDMode-CTRL-Action1-def CDMode-CTRL-Action2-def)
apply (*rule refl*)
done

12.2.7 Hierarchical Automaton *CarAudioSystem*

lemmas *CarAudioSystemStates = States-Root-CTRL States-CDPlayer-CTRL States-AudioPlayer-CTRL*
States-On-CTRL

States-TunerMode-CTRL States-CDMode-CTRL
Root-CTRL-States-def CDPlayer-CTRL-States-def
AudioPlayer-CTRL-States-def
On-CTRL-States-def TunerMode-CTRL-States-def
CDMode-CTRL-States-def

lemmas *CarAudioSystemInitState = Init-State-Root-CTRL Init-State-CDPlayer-CTRL*
Init-State-AudioPlayer-CTRL

Init-State-On-CTRL Init-State-TunerMode-CTRL
Init-State-CDMode-CTRL
Root-CTRL-Init-def CDPlayer-CTRL-Init-def
AudioPlayer-CTRL-Init-def
On-CTRL-Init-def TunerMode-CTRL-Init-def
CDMode-CTRL-Init-def

lemmas *CarAudioSystemEvents = Events-Root-CTRL Events-CDPlayer-CTRL*
Events-AudioPlayer-CTRL Events-On-CTRL

Events-TunerMode-CTRL Events-CDMode-CTRL

lemmas *CarAudioSystemthms = CarAudioSystemStates CarAudioSystemEvents*
CarAudioSystemInitState

schematic-goal *CarAudioSystem-StatesRoot*:

HASates (PseudoHA Root-CTRL (LiftInitData [V0 0, V1 0])) = ?X

apply (*wellformed CarAudioSystemthms*)+

done

lemmas *CarAudioSystemthms-1 = CarAudioSystemthms CarAudioSystem-StatesRoot*

schematic-goal *CarAudioSystem-StatesCDPlayer*:

HASates (PseudoHA Root-CTRL (LiftInitData [V0 0, V1 0]) [++]

("CarAudioSystem", CDPlayer-CTRL)) = ?X

apply (*wellformed CarAudioSystemthms-1*)+

done

lemmas *CarAudioSystemthms-2 = CarAudioSystemthms-1 CarAudioSystem-StatesCDPlayer*

schematic-goal *CarAudioSystem-StatesAudioPlayer*:

HASates (PseudoHA Root-CTRL (LiftInitData [V0 0, V1 0])

[++] ("CarAudioSystem", CDPlayer-CTRL)

[++] ("CarAudioSystem", AudioPlayer-CTRL)) = ?X

apply (*wellformed CarAudioSystemthms-2*)+

done

lemmas *CarAudioSystemthms-3 = CarAudioSystemthms-2 CarAudioSystem-StatesAudioPlayer*

schematic-goal *CarAudioSystem-StatesTunerMode*:

HASates (PseudoHA Root-CTRL (LiftInitData [V0 0, V1 0])

[++] ("CarAudioSystem", CDPlayer-CTRL)

[++] ("CarAudioSystem", AudioPlayer-CTRL)

[++] ("On", TunerMode-CTRL)) = ?X

apply (*wellformed CarAudioSystemthms-3*)+

done

lemmas *CarAudioSystemthms-4 = CarAudioSystemthms-3 CarAudioSystem-StatesTunerMode*

schematic-goal *CarAudioSystem-StatesCDMode*:

HASates (PseudoHA Root-CTRL (LiftInitData [V0 0, V1 0])

[++] ("CarAudioSystem", CDPlayer-CTRL)

[++] ("CarAudioSystem", AudioPlayer-CTRL)

[++] ("On", TunerMode-CTRL)

[++] ("On", CDMode-CTRL)) = ?X

apply (*wellformed CarAudioSystemthms-4*)+

done

lemmas *CarAudioSystemthms-5 = CarAudioSystemthms-4 CarAudioSystem-StatesCDMode*

schematic-goal *SAsCarAudioSystem*:

SAs CarAudioSystem = ?X

apply (*unfold CarAudioSystem-def*)

apply (*wellformed CarAudioSystemthms-5*)
done

schematic-goal *EventsCarAudioSystem*:
 HAEvents CarAudioSystem = ?X
apply (*unfold CarAudioSystem-def*)
apply (*wellformed CarAudioSystemthms-5*)
done

schematic-goal *CompFunCarAudioSystem*:
 CompFun CarAudioSystem = ?X
apply (*unfold CarAudioSystem-def*)
apply (*wellformed CarAudioSystemthms-5*)
done

schematic-goal *StatesCarAudioSystem*:
 HASates CarAudioSystem = ?X
apply (*unfold CarAudioSystem-def*)
apply (*wellformed CarAudioSystemthms-5*)
done

schematic-goal *ValueCarAudioSystem*:
 HAInitValue CarAudioSystem = ?X
apply (*unfold CarAudioSystem-def*)
apply (*wellformed CarAudioSystemthms-5*)
done

schematic-goal *HAInitStatesCarAudioSystem*:
 HAInitStates CarAudioSystem = ?X
by (*simp add: HAINitStates-def SAsCarAudioSystem CarAudioSystemInitState*)

schematic-goal *HARootCarAudioSystem*:
 HARoot CarAudioSystem = ?X
apply (*unfold CarAudioSystem-def*)
apply (*wellformed CarAudioSystemthms-5*)
done

schematic-goal *HAInitStateCarAudioSystem*:
 HAInitState CarAudioSystem = ?X
by (*simp add: HARootCarAudioSystem HAINitState-def CarAudioSystemInitState*)

lemma *check-DataSpace [simp]*:
 [*range V0, range V1*] \in *dataspace*
apply (*unfold dataspace-def DataSpace.DataSpace-def*)

```

apply auto
apply (rename-tac D)
apply (case-tac D)
apply auto
done

```

```

lemma PartNum-DataSpace [simp]:
   $PartNum\ DSspace = 2$ 
apply (unfold PartNum-def DSspace-def)
apply (simp add: Abs-dataspace-inverse)
done

```

```

lemma PartDom-DataSpace-V0 [simp]:
   $(PartDom\ DSspace\ 0) = range\ V0$ 
apply (unfold PartDom-def DSspace-def)
apply (simp add: Abs-dataspace-inverse)
done

```

```

lemma PartDom-DataSpace-V1 [simp]:
   $(PartDom\ DSspace\ (Suc\ 0)) = range\ V1$ 
apply (unfold PartDom-def DSspace-def)
apply (simp add: Abs-dataspace-inverse)
done

```

```

lemma check-InitialData [simp]:
   $([V0\ 0,\ V1\ 0], DSspace) \in data$ 
apply (unfold data-def Data.Data-def)
apply auto
apply (rename-tac d)
apply (case-tac d=0  $\vee$  d = 1)
apply auto
done

```

```

lemma Select0-InitData [simp]:
   $Select0\ (LiftInitData\ [V0\ 0,\ V1\ 0]) = 0$ 
apply (unfold LiftInitData-def Select0-def DataPart-def DataValue-def)
apply (simp add: Abs-data-inverse)
done

```

```

lemma Select1-InitData [simp]:
   $Select1\ (LiftInitData\ [V0\ 0,\ V1\ 0]) = 0$ 
apply (unfold LiftInitData-def Select1-def DataPart-def DataValue-def)
apply (simp add: Abs-data-inverse)
done

```

```

lemma HAINitValue1-CarAudioSystem:
   $CarAudioSystem\ |=H= Atom\ (VAL\ (\lambda\ d.\ (Select0\ d) = 0))$ 
apply (simp add: ValueCarAudioSystem)
done

```

```

lemma HAINitValue2-CarAudioSystem:
  CarAudioSystem |=H= Atom (VAL ( $\lambda d. (Select1\ d) = 0$ ))
apply (simp add: ValueCarAudioSystem)
done

lemma HAINitValue-DSpace-CarAudioSystem [simp]:
  Data.DataSpace (LiftInitData [V0 0, V1 0]) = DSpace
apply (unfold LiftInitData-def Data.DataSpace-def)
apply (simp add: Abs-data-inverse)
done

lemma check-InitStatus [simp]:
  (CarAudioSystem, InitConf CarAudioSystem, {},LiftInitData [V0 0, V1 0])  $\in$ 
  status
apply (unfold status-def Status-def)
apply (simp add: ValueCarAudioSystem)
done

lemma InitData-InitStatus [simp]:
  Value (InitStatus CarAudioSystem) = LiftInitData [V0 0, V1 0]
apply (simp add: ValueCarAudioSystem)
done

lemma Events-InitStatus [simp]:
  Events (InitStatus CarAudioSystem) = {}
apply (unfold InitStatus-def Events-def)
apply (simp add: Abs-status-inverse)
done

lemma Conf-InitStatus [simp]:
  Conf (InitStatus CarAudioSystem) = InitConf CarAudioSystem
apply (unfold InitStatus-def Conf-def)
apply (simp add: Abs-status-inverse)
done

lemma CompFunCarAudioSystem-the:
  the (CompFun CarAudioSystem "On") = {CDMode-CTRL, TunerMode-CTRL}
apply (unfold CarAudioSystem-def)
apply (subst AddSA-CompFun-the)
prefer 3
apply (subst AddSA-CompFun-the)
prefer 3
apply (subst AddSA-CompFun-the2)
apply (wellformed CarAudioSystemthms-5)+
done

lemma CompFunCarAudioSystem-the2:
  the (CompFun CarAudioSystem "CarAudioSystem") = {AudioPlayer-CTRL,

```

```

CDPlayer-CTRL}
apply (unfold CarAudioSystem-def)
apply (subst AddSA-CompFun-the3)
prefer 5
apply (subst AddSA-CompFun-the3)
prefer 5
apply (subst AddSA-CompFun-the)
prefer 3
apply (subst AddSA-CompFun-the)
prefer 3
apply (subst PseudoHA-CompFun-the)
apply (wellformed CarAudioSystemthms-5)+
done

```

```

lemma CompFunCarAudioSystem-the3:
  the (CompFun CarAudioSystem "Off") = {}
apply (unfold CarAudioSystem-def)
apply (subst AddSA-CompFun-the3)
prefer 5
apply (subst AddSA-CompFun-the3)
prefer 5
apply (subst AddSA-CompFun-the2)
apply (wellformed CarAudioSystemthms-5)+
done

```

```

schematic-goal CompFunCarAudioSystem-ran:
  ran (CompFun CarAudioSystem) = ?X
apply (unfold CarAudioSystem-def)
apply (rule AddSA-CompFun-ran3-IFF)
prefer 6
apply (subst AddSA-CompFun-PseudoHA-ran2)
prefer 4
apply (simp add: insert-commute)
apply (wellformed CarAudioSystemthms-5)+
done

```

```

lemma Root-CTRL-CDPlayer-CTRL-noteq [simp]:
  Root-CTRL ≠ CDPlayer-CTRL
apply (rule SA-States-disjunct)
apply (wellformed CarAudioSystemthms-5)+
done

```

```

lemma Root-CTRL-AudioPlayer-CTRL-noteq [simp]:
  Root-CTRL ≠ AudioPlayer-CTRL
apply (rule SA-States-disjunct)
apply (wellformed CarAudioSystemthms-5)+
done

```

```

lemma Root-CTRL-TunerMode-CTRL-noteq [simp]:

```


Root-CTRL \neq *TunerMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *Root-CTRL-CDMode-CTRL-noteq* [*simp*]:
Root-CTRL \neq *CDMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *CDPlayer-CTRL-AudioPlayer-CTRL-noteq* [*simp*]:
CDPlayer-CTRL \neq *AudioPlayer-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *CDPlayer-CTRL-TunerMode-CTRL-noteq* [*simp*]:
CDPlayer-CTRL \neq *TunerMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *CDPlayer-CTRL-CDMode-CTRL-noteq* [*simp*]:
CDPlayer-CTRL \neq *CDMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *AudioPlayer-CTRL-TunerMode-CTRL-noteq* [*simp*]:
AudioPlayer-CTRL \neq *TunerMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *AudioPlayer-CTRL-CDMode-CTRL-noteq* [*simp*]:
AudioPlayer-CTRL \neq *CDMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

lemma *TunerMode-CTRL-CDMode-CTRL-noteq* [*simp*]:
TunerMode-CTRL \neq *CDMode-CTRL*
apply (rule *SA-States-disjunct*)
apply (wellformed *CarAudioSystemthms-5*)
done

schematic-goal *Chi-CarAudioSystem*:
Chi CarAudioSystem "CarAudioSystem" = ?X

```

apply (unfold Chi-def)
apply (rule trans)
apply (simp add: SAsCarAudioSystem StatesCarAudioSystem restrict-def Comp-FunCarAudioSystem-the2)
apply (rule trans)
apply (simp add: not-sym)
apply (simp add: CarAudioSystemStates insert-or)
done

```

```

schematic-goal Chi-CarAudioSystem-On:
  Chi CarAudioSystem "On" = ?X
apply (unfold Chi-def)
apply (rule trans)
apply (simp add: SAsCarAudioSystem StatesCarAudioSystem restrict-def Comp-FunCarAudioSystem-the)
apply (rule trans)
apply (simp add: not-sym)
apply (simp add: CarAudioSystemStates insert-or)
done

```

```

schematic-goal Chi-CarAudioSystem-Off:
  Chi CarAudioSystem "Off" = ?X
apply (unfold Chi-def)
apply (simp add: SAsCarAudioSystem StatesCarAudioSystem restrict-def Comp-FunCarAudioSystem-the3)
done

```

```

schematic-goal InitConf-CarAudioSystem:
  InitConf CarAudioSystem = ?X
apply (unfold CarAudioSystem-def)
apply (rule AddSA-InitConf-IFF)+
apply simp
apply (wellformed CarAudioSystemthms-5)
apply fast
apply (wellformed CarAudioSystemthms-5)
apply simp
apply (simp add: CarAudioSystemthms-5)
apply (wellformed CarAudioSystemthms-5)
apply fast
apply (wellformed CarAudioSystemthms-5)
apply fast
apply simp
apply (wellformed CarAudioSystemthms-5)
apply fast
apply (wellformed CarAudioSystemthms-5)
apply fast
apply simp
apply (wellformed CarAudioSystemthms-5)
apply force

```

```
apply (wellformed CarAudioSystemthms-5)
apply fast
apply (simp add: CarAudioSystemthms-5)
done
```

```
lemma Initial-State-CarAudioSystem:
  CarAudioSystem |=H= Atom (IN "Off")
apply (simp add: InitConf-CarAudioSystem )
done
```

```
end
```

References

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