

Sophie Germain's Theorem

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1 Introduction

Fermat's Last Theorem (often abbreviated to FLT) states that for any integer $2 < n$, the equation $x^n + y^n = z^n$ has no nontrivial solution in the integers. Pierre de Fermat first conjectured this result in the 17th century, claiming to have a proof that did not fit in the margin of his notebook. However, it remained an open problem for centuries until Andrew Wiles

and Richard Taylor provided a complete proof in 1995 using advanced techniques from algebraic geometry and modular forms.

But in the meantime, many mathematicians have made partial progress on the problem. In particular, Sophie Germain's theorem states that p is a prime such that $2 * p + 1$ is also a prime, then there are no integer solutions to the equation $x^p + y^p = z^p$ such that p divides neither x , y nor z .

This result is not only included in the extended list of Freek's "Top 100 theorems"¹, but is also very familiar to students taking the French "agrégation" mathematics competitive examination. Hoping that this submission might also be useful to them, we developed separately the classical version of the theorem as presented in [1] and a generalization that one can find for example in [2].

¹<http://www.cs.ru.nl/~freek/100/>

The session displayed in 1 is organized as follows:

- `FLT_Sufficient_Conditions` provides sufficient conditions for proving FLT,
- `SG_Preliminaries` establish some useful lemmas and introduces the concept of Sophie Germain prime,
- `SG_Theorem` proves Sophie Germain's theorem and
- `SG_Generalization` gives a generalization of it.

2 Preliminaries

2.1 Coprimality

We start with this useful elimination rule: when a and b are not *coprime* and are not both equal to 0, there exists some common *prime* factor.

```
lemma (in factorial-semiring-gcd) not-coprime-nonzeroE :
  ⟨[¬ coprime a b; a ≠ 0 ∨ b ≠ 0; ∃ p. prime p ⇒ p dvd a ⇒ p dvd b ⇒
  thesis] ⇒ thesis⟩
  ⟨proof⟩
```

Still referring to the notion of *coprime* (but generalized to a set), we prove that when $\text{Gcd } A \neq 0$, the elements of $\{a \text{ div Gcd } A \mid a. a \in A\}$ are setwise *coprime*.

```
lemma (in semiring-Gcd) GCD-div-Gcd-is-one :
  ⟨(GCD a ∈ A. a div Gcd A) = 1⟩ if ⟨Gcd A ≠ 0⟩
  ⟨proof⟩
```

2.2 Power

Now we want to characterize the fact of admitting an n -th root with a condition on the *multiplicity* of each prime factor.

```
lemma exists-nth-root-iff :
  ⟨(∃ x. normalize y = x ^ n) ↔ (∀ p ∈ prime-factors y. n dvd multiplicity p y)⟩
  if ⟨y ≠ 0⟩ for y :: ⟨'a :: factorial-semiring-multiplicative⟩
  ⟨proof⟩
```

We use this result to obtain the following elimination rule.

```
corollary prod-is-some-powerE :
  fixes a b :: ⟨'a :: factorial-semiring-multiplicative⟩
```

```

assumes <coprime a b> and <a * b = x ^ n>
obtains α where <normalize a = α ^ n>
⟨proof⟩

```

2.3 Sophie Germain Prime

Finally, we introduce Sophie Germain primes.

```

definition SophGer-prime :: <nat ⇒ bool> (<SG>)
  where <SG p ≡ odd p ∧ prime p ∧ prime (2 * p + 1)>

```

```

lemma SophGer-primeI : <odd p ⇒ prime p ⇒ prime (2 * p + 1) ⇒ SG p>
  ⟨proof⟩

```

```

lemma SophGer-primeD : <odd p> <prime p> <prime (2 * p + 1)> if <SG p>
  ⟨proof⟩

```

We can easily compute Sophie Germain primes less than 2000.

```

value <[p. p ← [0..2000], SG (nat p)]>

```

```

context fixes p assumes <SG p> begin
  {ML}
  lemma nonzero : <p ≠ 0> ⟨proof⟩
  lemma pos : <0 < p> ⟨proof⟩
  lemma ge-3 : <3 ≤ p>
    ⟨proof⟩
  lemma ge-7 : <7 ≤ 2 * p + 1> ⟨proof⟩
  lemma notcong-zero :
    <[− 3 ≠ 0 :: int] (mod 2 * p + 1)> <[− 1 ≠ 0 :: int] (mod 2 * p + 1)>
    <[ 1 ≠ 0 :: int] (mod 2 * p + 1)> <[ 3 ≠ 0 :: int] (mod 2 * p + 1)>
    ⟨proof⟩
  lemma notcong-p :
    <[− 1 ≠ p :: int] (mod 2 * p + 1)>
    <[ 0 ≠ p :: int] (mod 2 * p + 1)>
    <[ 1 ≠ p :: int] (mod 2 * p + 1)>
    ⟨proof⟩
  lemma p-th-power-mod-q :
    <[m ^ p = 1] (mod 2 * p + 1) ∨ [m ^ p = − 1] (mod 2 * p + 1)>

```

```

if  $\neg 2 * p + 1 \text{ dvd } m$  for  $m :: \text{int}$ 
⟨proof⟩

```

```
end
```

2.4 Fermat's little Theorem for Integers

```

lemma fermat-theorem-int :
   $[a^{\wedge}(p - 1) = 1] \pmod{p}$  if prime  $p$  and  $\neg p \text{ dvd } a$ 
  for  $p :: \text{nat}$  and  $a :: \text{int}$ 
⟨proof⟩

```

3 Sufficient Conditions for FLT

Recall that FLT stands for “Fermat’s Last Theorem”. FLT states that there is no nontrivial integer solutions to the equation $x^n + y^n = z^n$ for any natural number $2 < n$. as soon as the natural number n is greater than 2. More formally: $2 < n \implies \nexists x y z. x^n + y^n = z^n$. We give here some sufficient conditions.

3.1 Coprimality

We first notice that it is sufficient to prove FLT for integers x , y and z that are (setwise) *coprime*.

```

lemma (in semiring-Gcd) FLT-setwise-coprime-reduction :
  assumes  $x^{\wedge}n + y^{\wedge}n = z^{\wedge}n$   $x \neq 0$   $y \neq 0$   $z \neq 0$ 
  defines  $d \equiv \text{Gcd}\{x, y, z\}$ 
  shows  $(x \text{ div } d)^{\wedge}n + (y \text{ div } d)^{\wedge}n = (z \text{ div } d)^{\wedge}n$   $x \text{ div } d \neq 0$ 
     $y \text{ div } d \neq 0$   $z \text{ div } d \neq 0$   $\text{Gcd}\{x \text{ div } d, y \text{ div } d, z \text{ div } d\} = 1$ 
⟨proof⟩

```

```

corollary (in semiring-Gcd) FLT-for-coprime-is-sufficient :
   $\nexists x y z. x \neq 0 \wedge y \neq 0 \wedge z \neq 0 \wedge \text{Gcd}\{x, y, z\} = 1 \wedge x^{\wedge}n + y^{\wedge}n = z^{\wedge}n$ 
 $\implies$ 
   $\nexists x y z. x \neq 0 \wedge y \neq 0 \wedge z \neq 0 \wedge x^{\wedge}n + y^{\wedge}n = z^{\wedge}n$ 
⟨proof⟩
lemma OFCLASS(int, semiring-Gcd-class) ⟨proof⟩

```

This version involving congruences will be useful later.

```

lemma FLT-setwise-coprime-reduction-mod-version :
  fixes  $x y z :: \text{int}$ 

```

```

assumes ⟨ $x \wedge n + y \wedge n = z \wedge n$ ⟩ ⟨ $x \neq 0$ ⟩ ⟨ $y \neq 0$ ⟩ ⟨ $z \neq 0$ ⟩ (mod m)
defines ⟨ $d \equiv Gcd\{x, y, z\}$ ⟩
shows ⟨ $(x \text{ div } d) \wedge n + (y \text{ div } d) \wedge n = (z \text{ div } d) \wedge n$ ⟩ ⟨ $x \text{ div } d \neq 0$ ⟩ (mod m)
      ⟨ $y \text{ div } d \neq 0$ ⟩ (mod m) ⟨ $z \text{ div } d \neq 0$ ⟩ (mod m) ⟨ $Gcd\{x \text{ div } d, y \text{ div } d, z \text{ div } d\} = 1$ ⟩
⟨proof⟩

```

Actually, it is sufficient to prove FLT for integers x, y and z that are pairwise *coprime*

```

lemma (in semiring-Gcd) FLT-setwise-coprime-imp-pairwise-coprime :
  ⟨coprime x y⟩ if ⟨ $n \neq 0$ ⟩ ⟨ $x \wedge n + y \wedge n = z \wedge n$ ⟩ ⟨ $Gcd\{x, y, z\} = 1$ ⟩
⟨proof⟩

```

3.2 Odd prime Exponents

From Fermat3_4, FLT is already proven for $n = 4$. Using this, we can prove that it is sufficient to prove FLT for *odd prime* exponents.

```

lemma (in semiring-1-no-zero-divisors) FLT-exponent-reduction :
  assumes ⟨ $x \wedge n + y \wedge n = z \wedge n$ ⟩ ⟨ $x \neq 0$ ⟩ ⟨ $y \neq 0$ ⟩ ⟨ $z \neq 0$ ⟩ ⟨ $p \text{ dvd } n$ ⟩
  shows ⟨ $(x \wedge (n \text{ div } p)) \wedge p + (y \wedge (n \text{ div } p)) \wedge p = (z \wedge (n \text{ div } p)) \wedge p$ ⟩
      ⟨ $x \wedge (n \text{ div } p) \neq 0$ ⟩ ⟨ $y \wedge (n \text{ div } p) \neq 0$ ⟩ ⟨ $z \wedge (n \text{ div } p) \neq 0$ ⟩
⟨proof⟩

```

```
lemma ⟨OFCLASS(int, semiring-1-no-zero-divisors-class)⟩ ⟨proof⟩
```

```

lemma odd-prime-or-four-factorE :
  fixes  $n :: nat$  assumes ⟨ $2 < n$ ⟩
  obtains  $p$  where ⟨ $p \text{ dvd } n$ ⟩ ⟨ $odd\ p$ ⟩ ⟨ $prime\ p$ ⟩ | ⟨ $4 \text{ dvd } n$ ⟩
⟨proof⟩

```

Finally, proving FLT for odd prime exponents is sufficient.

```

corollary FLT-for-odd-prime-exponents-is-sufficient :
  ⟨ $\nexists x y z :: int. x \neq 0 \wedge y \neq 0 \wedge z \neq 0 \wedge x \wedge n + y \wedge n = z \wedge n$ ⟩ if ⟨ $2 < n$ ⟩
  and odd-prime-FLT :
    ⟨ $\bigwedge p. odd\ p \implies prime\ p \implies$ 
      ⟨ $\nexists x y z :: int. x \neq 0 \wedge y \neq 0 \wedge z \neq 0 \wedge x \wedge p + y \wedge p = z \wedge p$ ⟩
⟨proof⟩

```

4 Sophie Germain's Theorem: classical Version

The proof we give here is from [1].

4.1 A Crucial Lemma

```

lemma Sophie-Germain-lemma-computation :
  fixes x y :: int assumes <odd p>
  defines <S ≡ ∑ k = 0..p - 1. (-y)^(p-1-k) * x^k>
  shows <(x + y) * S = x^p + y^p>
  <proof>

lemma Sophie-Germain-lemma-computation-cong-simp :
  fixes p :: nat and n x y :: int assumes <p ≠ 0> <[y = -x] (mod n)>
  defines <S ≡ λx y. ∑ k = 0..p - 1. (-y)^(p-1-k) * x^k>
  shows <[S x y = p * x^(p-1)] (mod n)>
  <proof>

lemma Sophie-Germain-lemma-only-possible-prime-common-divisor :
  fixes x y z :: int and p :: nat
  defines S-def: <S ≡ λx y. ∑ k = 0..p - 1. (-y)^(p-1-k) * x^k>
  assumes <prime p> <prime q> <coprime x y> <q dvd x + y> <q dvd S x y>
  shows <q = p>
  <proof>

lemma Sophie-Germain-lemma :
  fixes x y z :: int
  assumes <odd p> and <prime p> and fermat : <x^p + y^p + z^p = 0>
    and <[x ≠ 0] (mod p)> and <coprime y z>
  defines <S ≡ ∑ k = 0..p - 1. (-z)^(p-1-k) * y^k>
  shows <∃ a α. y + z = a^p ∧ S = α^p>
  <proof>
```

4.2 The Theorem

```

theorem Sophie-Germain-theorem :
  <¬∃ x y z :: int. x^p + y^p = z^p ∧ [x ≠ 0] (mod p) ∧
    [y ≠ 0] (mod p) ∧ [z ≠ 0] (mod p)> if SG : <SG p>
  <proof>
```

5 Sophie Germain's Theorem: generalized Version

The proof we give here is from [2].

5.1 Auxiliary Primes

```

abbreviation non-consecutivity-condition :: <nat ⇒ nat ⇒ bool> (<NC>)
  where <NC p q ≡ ¬∃ x y :: int. [x ≠ 0] (mod q) ∧ [y ≠ 0] (mod q) ∧ [x^p = 1
    + y^p] (mod q)>
```

lemma *non-consecutivity-condition-bis* :
 $\langle NC\ p\ q \longleftrightarrow (\nexists x\ y\ a\ b. [a :: int \neq 0] (mod q) \wedge [a \wedge p = x] (mod q) \wedge [b :: int \neq 0] (mod q) \wedge [b \wedge p = y] (mod q) \wedge [x = 1 + y] (mod q)) \rangle$
 $\langle proof \rangle$

abbreviation *not-pth-power* :: $\langle nat \Rightarrow nat \Rightarrow bool \rangle$ ($\langle PPP \rangle$)
where $\langle PPP\ p\ q \equiv \nexists x :: int. [p = x \wedge p] (mod q) \rangle$

definition *auxiliary-prime* :: $\langle nat \Rightarrow nat \Rightarrow bool \rangle$ ($\langle aux'-prime \rangle$)
where $\langle aux-prime\ p\ q \equiv prime\ p \wedge prime\ q \wedge NC\ p\ q \wedge PPP\ p\ q \rangle$

lemma *auxiliary-primeI* :
 $\langle [\![prime\ p; prime\ q; NC\ p\ q; PPP\ p\ q]\!] \implies aux-prime\ p\ q \rangle$
 $\langle proof \rangle$

lemma *auxiliary-primeD* :
 $\langle prime\ p \wedge prime\ q \wedge NC\ p\ q \wedge PPP\ p\ q \text{ if } aux-prime\ p\ q \rangle$
 $\langle proof \rangle$

We do not give code equation yet, let us first work around these notions.

lemma *gen-mult-group-mod-prime-as-ord* : $\langle ord\ p\ g = p - 1 \rangle$
if $\langle prime\ p \rangle \langle \{1 .. p - 1\} = \{g \wedge k \text{ mod } p | k \in UNIV\} \rangle$
 $\langle proof \rangle$

lemma *exists-nth-power-mod-prime-iff* :
fixes $p\ n$ **assumes** $\langle prime\ p \rangle$
defines $d\text{-def} : \langle d \equiv gcd\ n\ (p - 1) \rangle$
shows $\langle (\exists x :: int. [a = x \wedge n] (mod p)) \longleftrightarrow (n \neq 0 \wedge [a = 0] (mod p)) \vee [a \wedge ((p - 1) \text{ div } d) = 1] (mod p) \rangle$
 $\langle proof \rangle$

corollary *not-pth-power-iff* :
 $\langle PPP\ p\ q \longleftrightarrow [p \neq 0] (mod q) \wedge [p \wedge ((q - 1) \text{ div } gcd\ p\ (q - 1)) \neq 1] (mod q) \rangle$
if $\langle prime\ p \rangle \langle prime\ q \rangle$
 $\langle proof \rangle$

corollary *not-pth-power-iff-mod* :
 $\langle PPP\ p\ q \longleftrightarrow \neg q \text{ dvd } p \wedge p \wedge ((q - 1) \text{ div } gcd\ p\ (q - 1)) \text{ mod } q \neq 1 \rangle$
if $\langle prime\ p \rangle$ **and** $\langle prime\ q \rangle$
 $\langle proof \rangle$

lemma *non-consecutivity-condition-iff-enum-mod* :

— This version is oriented towards code generation.

```
<NC p q <→
  (forall x in {1..q - 1}. let x-p-mod = x ^ p mod q
   in forall y in {1..q - 1}. x-p-mod ≠ (1 + y ^ p mod q) mod q)
  if ⟨q ≠ 0⟩
⟨proof⟩
```

lemma auxiliary-prime-iff-enum-mod [code] :

— We will have a more optimized version later.

```
<aux-prime p q <→
  prime p ∧ prime q ∧
  ¬ q dvd p ∧ p ^ ((q - 1) div gcd p (q - 1)) mod q ≠ 1 ∧
  (forall x in {1..q - 1}. let x-p-mod = x ^ p mod q
   in forall y in {1..q - 1}. x-p-mod ≠ (1 + y ^ p mod q) mod q)
⟨proof⟩
```

We can for example compute pairs of auxiliary primes less than 110.

```
value <[(p, q). p ← [1..110], q ← [1..110], aux-prime (nat p) (nat q)]>
```

lemma auxiliary-primeI' :

```
<[prime p; prime q; ¬ q dvd p; p ^ ((q - 1) div gcd p (q - 1)) mod q ≠ 1;
  ∧ x y. x ∈ {1..q - 1} ⇒ y ∈ {1..q - 1} ⇒ [x ^ p ≠ 1 + y ^ p] (mod q)]>
  ⇒ aux-prime p q
⟨proof⟩
```

lemma two-is-not-auxiliary-prime : <¬ aux-prime p 2>

⟨proof⟩

lemma auxiliary-prime-of-2 : <aux-prime 2 q <→ q = 3 ∨ q = 5>

An auxiliary prime q of p is generally of the form $q = (2::'a) * n * p + 1$.

lemma auxiliary-prime-pattern-aux :

```
<exists x y. [x ≠ 0] (mod q) ∧ [y ≠ 0] (mod q) ∧ [x ^ p = 1 + y ^ p] (mod q)
  if ⟨p ≠ 0⟩ ⟨prime q⟩ ⟨coprime p (q - 1)⟩ ⟨odd q⟩
⟨proof⟩
```

theorem auxiliary-prime-pattern :

```
<p = 2 ∧ (q = 3 ∨ q = 5) ∨ odd p ∧ (exists n ≥ 1. q = 2 * n * p + 1)> if aux-p :
  <aux-prime p q>
⟨proof⟩
```

```
lemma auxiliary-prime-imp-less : <aux-prime p q  $\implies$  p < q>
  ⟨proof⟩
```

```
lemma auxiliary-primeE :
  assumes <aux-prime p q>
  obtains <p = 2> <q = 3> | <p = 2> <q = 5>
    | n where <odd p> <1 ≤ n> <q = 2 * n * p + 1>
      <NC p (2 * n * p + 1)> <PPP p (2 * n * p + 1)>
  ⟨proof⟩
```

With this, we can quickly eliminate numbers that cannot be auxiliary.

```
declare auxiliary-prime-iff-enum-mod [code del]
```

```
lemma auxiliary-prime-iff-enum-mod-optimized [code] :
  <aux-prime p q  $\longleftrightarrow$ 
    p = 2  $\wedge$  (q = 3  $\vee$  q = 5)  $\vee$ 
    p < q  $\wedge$ 
    2 * p  $\text{dvd}$  q - 1  $\wedge$ 
    prime p  $\wedge$  prime q  $\wedge$ 
     $\neg$  q  $\text{dvd}$  p  $\wedge$  p  $\hat{\wedge}$ ((q - 1)  $\text{div gcd}$  p (q - 1))  $\text{mod}$  q  $\neq$  1  $\wedge$ 
    ( $\forall x \in \{1..q - 1\}$ . let x-p-mod = x  $\hat{\wedge}$  p  $\text{mod}$  q
      in  $\forall y \in \{1..q - 1\}$ . x-p-mod  $\neq$  (1 + y  $\hat{\wedge}$  p  $\text{mod}$  q)  $\text{mod}$  q)>
  ⟨proof⟩
```

```
value <[(p, q). p  $\leftarrow$  [1..1000], q  $\leftarrow$  [1..110], aux-prime (nat p) (nat q)]>
```

5.2 Sophie Germain Primes are auxiliary

When p is an *odd prime* and $2 * p + 1$ is also a *prime* (what we call a *Sophie Germain prime*), $2 * p + 1$ is automatically an *aux-prime*.

```
lemma SophGer-prime-imp-auxiliary-prime :
  fixes p assumes <SG p> defines q-def: <q  $\equiv$  2 * p + 1>
  shows <aux-prime p q>
  ⟨proof⟩
```

5.3 Main Theorems

```
theorem Sophie-Germain-auxiliary-prime :
  <q  $\text{dvd}$  x  $\vee$  q  $\text{dvd}$  y  $\vee$  q  $\text{dvd}$  z>
  if <x  $\hat{\wedge}$  p + y  $\hat{\wedge}$  p = z  $\hat{\wedge}$  p> and <aux-prime p q> for x y z :: int
  ⟨proof⟩
```

```
theorem Sophie-Germain-generalization :
   $\nexists x y z :: \text{int}. x \hat{\wedge} p + y \hat{\wedge} p = z \hat{\wedge} p \wedge$ 
    [x  $\neq 0$ ] (mod  $p^2$ )  $\wedge$  [y  $\neq 0$ ] (mod  $p^2$ )  $\wedge$  [z  $\neq 0$ ] (mod  $p^2$ )
  if odd-p : <odd p> and aux-prime : <aux-prime p q>
```

$\langle proof \rangle$

Since $SG\ p \implies aux\text{-}prime\ p\ (2 * p + 1)$, this result is a generalization of $SG\ p \implies \nexists x\ y\ z.\ x^p + y^p = z^p \wedge [x \neq 0] \pmod{int\ p} \wedge [y \neq 0] \pmod{int\ p} \wedge [z \neq 0] \pmod{int\ p}$.

References

- [1] S. Francinou, H. Gianella, and S. Nicolas. *Oraux X-ENS Algèbre 1*. Cassini, 2014.
- [2] A. Kiefer. Le théorème de Fermat vu par M. Le Blanc. *Brussels Summer School of Mathematics, Notes de la cinquième BSSM*, 2012.

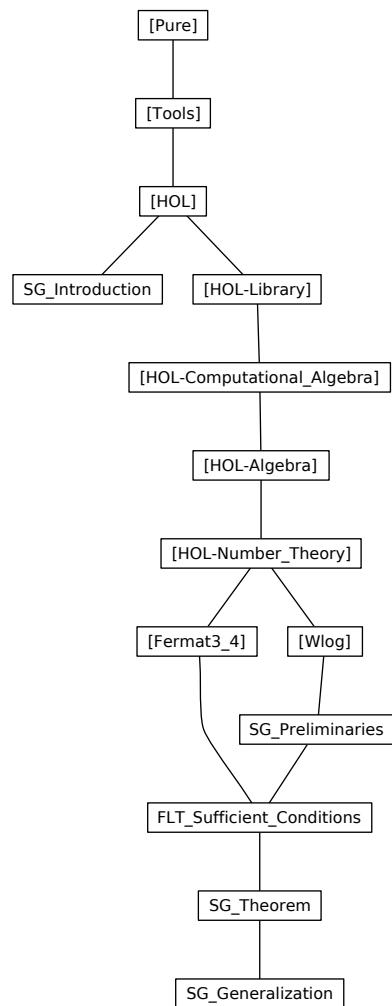


Figure 1: Dependency graph of the session `Sophie_Germain`