Exploring Simplified Variants of Gödel's Ontological Argument in Isabelle/HOL

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Abstract

Simplified variants of Gödel's ontological argument are explored. Among those is a particularly interesting simplified argument which is (i) valid already in basic modal logics K or KT, (ii) which does not suffer from modal collapse, and (iii) which avoids the rather complex predicates of essence (Ess.) and necessary existence (NE) as used by Gödel.

Whether the presented variants increase or decrease the attractiveness and persuasiveness of the ontological argument is a question I would like to pass on to philosophy and theology.

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1 Background Reading

The selected simplified variants of Gödel's ontological argument [5, 7] as presented in §3 have first been extracted from the insights gained in §6 of [4]. These variants are also influenced by the work presented in [1] and they significantly extend the findings from [3]. In §4 we additionally include the sources from [4].

2 Higher-Order Modal Logic in HOL (cf. [2] and Fig. 1 in [4]).

```
theory HOML imports Main begin nitpick-params[user-axioms, expect=genuine]
```

Type i is associated with possible worlds and type e with entities:

```
typedecl i — Possible worlds typedecl e — Individuals type-synonym \sigma = i \Rightarrow bool — World-lifted propositions type-synonym \gamma = e \Rightarrow \sigma — Lifted predicates type-synonym \mu = \sigma \Rightarrow \sigma — Unary modal connectives type-synonym \nu = \sigma \Rightarrow \sigma \Rightarrow \sigma — Binary modal connectives
```

Logical connectives (operating on truth-sets):

```
abbreviation c1::\sigma (\langle\bot\rangle) where \bot \equiv \lambda w. False abbreviation c2::\sigma (\langle\top\rangle) where \top \equiv \lambda w. True abbreviation c3::\mu (\langle\neg\neg\cdot|52|53\rangle) where \neg\varphi \equiv \lambda w.\neg(\varphi\ w) abbreviation c4::\nu (infix \langle\land\rangle 50\rangle) where \varphi\land\psi \equiv \lambda w.(\varphi\ w)\land(\psi\ w) abbreviation c5::\nu (infix \langle\rightarrow\rangle 49\rangle) where \varphi\lor\psi \equiv \lambda w.(\varphi\ w)\lor(\psi\ w) abbreviation c6::\nu (infix \langle\rightarrow\rangle 48\rangle) where \varphi\to\psi \equiv \lambda w.(\varphi\ w)\to(\psi\ w) abbreviation c7::\nu (infix \langle\leftrightarrow\rangle 47\rangle) where \varphi\leftrightarrow\psi \equiv \lambda w.(\varphi\ w)\longleftrightarrow(\psi\ w) consts R::i\Rightarrow i\Rightarrow bool (\langle\neg\neg\rangle) — Accessibility relation abbreviation c8::\mu (\langle\Box\cdot|54|55\rangle) where \Box\varphi \equiv \lambda w.\forall\ v.(wrv)\to(\varphi\ v) abbreviation c9::\mu (\langle\diamondsuit\neg\cdot|54|55\rangle) where \Diamond\varphi \equiv \lambda w.\exists\ v.(wrv)\land(\varphi\ v) abbreviation c10::\gamma\Rightarrow\gamma (\langle\neg\neg\rangle) where \neg\Phi \equiv \lambda x.\lambda w.\neg(\Phi\ x\ w) abbreviation c11::\gamma\Rightarrow\gamma (\langle\neg\neg\rangle) where \neg\Phi \equiv \lambda x.\lambda w.\neg(\Phi\ x\ w) abbreviation c12::e\Rightarrow e\Rightarrow\sigma (\langle\neg\neg\rangle) where x=y\equiv\lambda w.(x=y) abbreviation c13::e\Rightarrow e\Rightarrow\sigma (\langle\neg\neg\rightarrow\rangle) where x\neq y\equiv\lambda w.(x\neq y)
```

Polymorphic possibilist quantification:

```
abbreviation q1::('a\Rightarrow\sigma)\Rightarrow\sigma (\langle\forall\;\rangle) where \forall\;\Phi\equiv\lambda w.\forall\;x.(\Phi\;x\;w) abbreviation q2 (binder\langle\forall\;\rangle[10]11) where \forall\;x.\;\varphi(x)\equiv\forall\;\varphi abbreviation q3::('a\Rightarrow\sigma)\Rightarrow\sigma (\langle\exists\;\rangle) where \exists\;\Phi\equiv\lambda w.\exists\;x.(\Phi\;x\;w) abbreviation q4 (binder\langle\exists\;\rangle[10]11) where \exists\;x.\;\varphi(x)\equiv\exists\;\varphi
```

Actualist quantification for individuals/entities:

```
consts existsAt::\gamma (<-@->) abbreviation q5::\gamma\Rightarrow\sigma (\forall ^E \Rightarrow) where \forall ^E \Phi \equiv \lambda w. \forall x. (x@w) \longrightarrow (\Phi \ x \ w) abbreviation q6 (binder(\forall ^E \Rightarrow [8]9)) where \forall ^E x. \ \varphi(x) \equiv \forall ^E \varphi abbreviation q7::\gamma\Rightarrow\sigma ((\exists ^E \Rightarrow)) where \exists ^E \Phi \equiv \lambda w. \exists x. (x@w) \land (\Phi \ x \ w) abbreviation q8 (binder(\exists ^E \Rightarrow [8]9)) where \exists ^E x. \ \varphi(x) \equiv \exists ^E \varphi
```

Meta-logical predicate for global validity:

```
abbreviation g1::\sigma\Rightarrow bool\ (\langle |-| \rangle) where |\psi| \equiv \forall w. \psi w
```

Barcan and converse Barcan formula:

```
lemma True\ \mathbf{nitpick}[satisfy]\ \langle proof \rangle

lemma \lfloor (\forall^E x. \Box(\varphi\ x)) \to \Box(\forall^E x. (\varphi\ x)) \rfloor\ \mathbf{nitpick}\ \langle proof \rangle

lemma \lfloor \Box(\forall^E x. (\varphi\ x)) \to (\forall^E x. \Box(\varphi\ x)) \rfloor\ \mathbf{nitpick}\ \langle proof \rangle

lemma \lfloor (\forall\ x. \Box(\varphi\ x)) \to \Box(\forall\ x.\ \varphi\ x) \rfloor\ \langle proof \rangle

lemma \lfloor \Box(\forall\ x. (\varphi\ x)) \to (\forall\ x. \Box(\varphi\ x)) \rfloor\ \langle proof \rangle

end
```

3 Selected Simplified Ontological Argument

```
theory SimplifiedOntologicalArgument imports
HOML
begin
```

Positive properties:

```
consts posProp::\gamma \Rightarrow \sigma (\langle \mathcal{P} \rangle)
```

An entity x is God-like if it possesses all positive properties.

```
definition G(\langle \mathcal{G} \rangle) where \mathcal{G}(x) \equiv \forall \Phi.(\mathcal{P}(\Phi) \to \Phi(x))
```

The axiom's of the simplified variant are presented next; these axioms are further motivated in [4, 1]).

Self-difference is not a positive property (possible alternative: the empty property is not a positive property).

```
axiomatization where CORO1: |\neg(\mathcal{P}(\lambda x.(x\neq x)))|
```

A property entailed by a positive property is positive.

```
axiomatization where CORO2: |\forall \Phi \Psi. \mathcal{P}(\Phi) \land (\forall x. \Phi(x) \to \Psi(x)) \to \mathcal{P}(\Psi)|
```

Being Godlike is a positive property.

axiomatization where AXIOM3: |P G|

3.1 Verifying the Selected Simplified Ontological Argument (version 1)

The existence of a non-exemplified positive property implies that self-difference (or, alternatively, the empty property) is a positive property.

```
lemma LEMMA1: \lfloor (\exists \Phi. (\mathcal{P}(\Phi) \land \neg(\exists x. \Phi(x)))) \rightarrow \mathcal{P}(\lambda x. (x \neq x)) \rfloor \land proof \rangle
```

A non-exemplified positive property does not exist.

```
lemma LEMMA2: [\neg(\exists \Phi.(\mathcal{P}(\Phi) \land \neg(\exists x. \Phi(x))))] \langle proof \rangle
```

Positive properties are exemplified.

```
lemma LEMMA3: [\forall \Phi.(\mathcal{P}(\Phi) \rightarrow (\exists x. \Phi(x)))] \langle proof \rangle
```

There exists a God-like entity.

```
theorem THEOREM3': [\exists x. \mathcal{G}(x)] \langle proof \rangle
```

Necessarily, there exists a God-like entity.

```
theorem THEOREM3: [\Box(\exists x. \mathcal{G}(x))] \langle proof \rangle
```

However, the possible existence of Godlike entity is not implied.

```
theorem CORO: [\diamondsuit(\exists x. \mathcal{G}(x))] nitpick \langle proof \rangle
```

3.2 Verifying the Selected Simplified Ontological Argument (version 2)

We switch to logic T.

```
\begin{array}{ll} \textbf{axiomatization where} \ T \colon [\forall \, \varphi. \ \Box \varphi \to \varphi] \\ \textbf{lemma} \ T' \colon [\forall \, \varphi. \ \varphi \to \diamond \varphi] \ \langle \textit{proof} \rangle \end{array}
```

Positive properties are possibly exemplified.

```
theorem THEOREM1: [\forall \Phi. \mathcal{P}(\Phi) \rightarrow \Diamond(\exists x. \Phi(x))] \langle proof \rangle
```

Possibly there exists a God-like entity.

```
theorem CORO: [\diamondsuit(\exists x. \mathcal{G}(x))] \langle proof \rangle
```

The possible existence of a God-like entity impplies the necessary existence of a God-like entity.

```
theorem THEOREM2: [\diamondsuit(\exists x. \mathcal{G}(x)) \rightarrow \Box(\exists x. \mathcal{G}(x))] \land proof \rangle
```

```
Necessarily, there exists a God-like entity.
```

```
theorem THEO3: [\Box(\exists x. \mathcal{G}(x))] \langle proof \rangle
```

There exists a God-like entity.

```
theorem THEO3': [\exists x. \mathcal{G}(x)] \langle proof \rangle
```

Modal collapse is not implied; nitpick reports a countermodel.

```
lemma MC: [\forall \Phi. \Phi \rightarrow \Box \Phi] \text{ nitpick } \langle proof \rangle
```

Consistency of the theory; nitpick reports a model.

lemma $True \ \mathbf{nitpick}[\mathit{satisfy}] \ \langle \mathit{proof} \rangle$ end

4 Presentation of All Variants as Studied in [4]

4.1 Preliminaries: Modal Ultrafilter (Fig. 2 in [4])

theory MFilter imports HOML begin

Some abbreviations for auxiliary operations.

```
abbreviation a::\gamma\Rightarrow(\gamma\Rightarrow\sigma)\Rightarrow\sigma (\leftarrow\leftarrow\rightarrow) where x\in S\equiv S x abbreviation b::\gamma (\leftarrow0) where \emptyset\equiv\lambda x. \bot abbreviation c::\gamma (\leftarrow1) where U\equiv\lambda x. \top abbreviation d::\gamma\Rightarrow\gamma\Rightarrow\sigma (\leftarrow\leftarrow\leftarrow\rightarrow2) where \varphi\subseteq\psi\equiv\forall x.((\varphi\ x)\to(\psi\ x)) abbreviation e::\gamma\Rightarrow\gamma\Rightarrow\gamma (\leftarrow\leftarrow\leftarrow\rightarrow\rightarrow2) where \varphi\sqcap\psi\equiv\lambda x.((\varphi\ x)\land(\psi\ x)) abbreviation f::\gamma\Rightarrow\gamma (\leftarrow\leftarrow\leftarrow\rightarrow\rightarrow2) where \Rightarrow1 \Rightarrow2 \Rightarrow3 \Rightarrow4.
```

Definition of modal filter.

```
abbreviation g::(\gamma\Rightarrow\sigma)\Rightarrow\sigma\ (\langle Filter\rangle) where Filter\ \Phi\equiv(((\mathbf{U}\in\Phi)\ \land\ \neg(\emptyset\in\Phi)) \land\ (\forall\ \varphi\ \psi.\ (((\varphi\in\Phi)\ \land\ (\varphi\subseteq\psi))\ \rightarrow\ (\psi\in\Phi)))) \land\ (\forall\ \varphi\ \psi.\ (((\varphi\in\Phi)\ \land\ (\psi\in\Phi))\ \rightarrow\ ((\varphi\sqcap\psi)\in\Phi)))
```

Definition of modal ultrafilter .

```
abbreviation h::(\gamma \Rightarrow \sigma) \Rightarrow \sigma \ (\langle \mathit{UFilter} \rangle) where \mathit{UFilter} \ \Phi \equiv (\mathit{Filter} \ \Phi) \land (\forall \varphi.((\varphi \in \Phi) \lor ((^{-1}\varphi) \in \Phi)))
```

Modal filter and modal ultrafilter are consistent.

```
lemma [\forall \Phi \varphi.((\mathit{UFilter} \Phi) \to \neg((\varphi \in \Phi) \land ((^{-1}\varphi) \in \Phi)))] \langle \mathit{proof} \rangle  end
```

4.2 Preliminaries: Further Basic Notions (Fig. 3 in [4])

theory BaseDefs imports HOML begin

```
Positive properties.
consts posProp::\gamma \Rightarrow \sigma (\langle \mathcal{P} \rangle)
     Basic definitions for modal ontological argument.
abbreviation a (\langle - \square - \rangle) where X \square Y \equiv \forall E z.((X z) \rightarrow (Y z))
abbreviation b (\langle - \Rightarrow - \rangle) where X \Rightarrow Y \equiv \Box(X \sqsubseteq Y)
abbreviation c (\langle \mathcal{P} os \rangle) where \mathcal{P} os Z \equiv \forall X.((Z X) \rightarrow (\mathcal{P} X))
u))))))
     Definition of Godlike.
definition G(\langle \mathcal{G} \rangle) where \mathcal{G} x \equiv \forall Y.((\mathcal{P} Y) \rightarrow (Y x))
     Definitions of Essence and Necessary Existence.
definition E(\langle \mathcal{E} \rangle) where \mathcal{E}(Yx) = (Yx) \land (\forall Z.((Zx) \rightarrow (Y \Rightarrow Z)))
definition NE (\langle \mathcal{NE} \rangle) where \mathcal{NE} x \equiv \forall Y . ((\mathcal{E} Y x) \rightarrow \Box(\exists^E Y))
end
4.3
          Ultrafilter Analysis of Scott's Variant (Fig. 3 in [4]))
theory Scott Variant imports
  HOML
  MFilter
  BaseDefs
begin
     Axioms of Scott's variant.
axiomatization where
  A1: |\forall X.((\neg(\mathcal{P} X)) \leftrightarrow (\mathcal{P}(\rightarrow X)))| and
  A2: |\forall X \ Y.(((\mathcal{P} \ X) \land (X \Longrightarrow Y)) \rightarrow (\mathcal{P} \ Y))| and
  A3: |\forall \mathcal{Z}.((\mathcal{P}os \mathcal{Z}) \to (\forall X.((X \square \mathcal{Z}) \to (\mathcal{P}X))))| and
  A_4: |\forall X.((\mathcal{P} X) \rightarrow \Box(\mathcal{P} X))| and
  A5: |\mathcal{P}| \mathcal{NE}| and
  B: |\forall \varphi.(\varphi \to \Box \Diamond \varphi)| — Logic KB
lemma B': \forall x \ y. \ \neg(x\mathbf{r}y) \lor (y\mathbf{r}x) \ \langle proof \rangle
     Necessary existence of a Godlike entity.
theorem T6: \lfloor \Box(\exists^E \mathcal{G}) \rfloor
\langle proof \rangle
     Existence of a Godlike entity.
lemma |\exists^E \mathcal{G}| \langle proof \rangle
     Consistency
lemma True nitpick[satisfy] \langle proof \rangle
     Modal collapse: holds.
```

lemma $MC: |\forall \Phi.(\Phi \rightarrow \Box \Phi)|$

```
\langle proof \rangle
     Analysis of positive properties using ultrafilters.
theorem U1: |UFilter P| — Proof found by sledgehammer
\langle proof \rangle
lemma L1: |\forall X \ Y.((X \Rrightarrow Y) \to (X \sqsubseteq Y))| \langle proof \rangle
lemma L2: |\forall X \ Y.(((\mathcal{P} \ X) \land (X \sqsubseteq Y)) \rightarrow (\mathcal{P} \ Y))| \ \langle proof \rangle
     Set of supersets of X, we call this HF X.
abbreviation HF where HF X \equiv \lambda Y . (X \sqsubseteq Y)
     HF \mathcal{G} is a filter; hence, HF \mathcal{G} is Hauptfilter of \mathcal{G}.
lemma F1: |Filter (HF \mathcal{G})| \langle proof \rangle
lemma F2: |UFilter(HF \mathcal{G})| \langle proof \rangle
     T6 follows directly from F1.
theorem T6again: |\Box(\exists^E \mathcal{G})| \langle proof \rangle
end
         Ultrafilter Variant (Fig. 5 in [4])
4.4
theory UFilterVariant imports
  HOML
  MFilter
  BaseDefs
begin
     Axiom's of ultrafilter variant.
axiomatization where
  U1: |UFilter \mathcal{P}| and
  A2: |\forall X \ Y.(((\mathcal{P} \ X) \land (X \Longrightarrow Y)) \rightarrow (\mathcal{P} \ Y))| and
  A3: \left[ \forall \, \mathcal{Z}.((\mathcal{P}\text{os }\mathcal{Z}) \to (\forall \, X.((X \square \mathcal{Z}) \to (\mathcal{P} \, X)))) \right]
     Necessary existence of a Godlike entity.
theorem T6: \lfloor \Box(\exists^E \mathcal{G}) \rfloor — Proof also found by sledgehammer
\langle proof \rangle
     Checking for consistency.
lemma True nitpick[satisfy] \langle proof \rangle
     Checking for modal collapse.
lemma MC: |\forall \Phi.(\Phi \rightarrow \Box \Phi)| \text{ nitpick } \langle proof \rangle
```

```
4.5 Simplified Variant (Fig. 6 in [4])
```

```
theory SimpleVariant imports
  HOML
  MFilter
  BaseDefs
begin
     Axiom's of new, simplified variant.
axiomatization where
  A1': |\neg(\mathcal{P}(\lambda x.(x\neq x)))| and
  A2': |\forall X \ Y.(((\mathcal{P} \ X) \land ((X \sqsubseteq Y) \lor (X \Rrightarrow Y))) \rightarrow (\mathcal{P} \ Y))| and
  A3: |\forall \mathcal{Z}.((\mathcal{P}os \mathcal{Z}) \to (\forall X.((X \square \mathcal{Z}) \to (\mathcal{P} X))))|
lemma T2: |\mathcal{P}| \mathcal{G}| \langle proof \rangle
lemma L1: |\mathcal{P}(\lambda x.(x=x))| \langle proof \rangle
     Necessary existence of a Godlike entity.
theorem T6: |\Box(\exists^E \mathcal{G})| — Proof found by sledgehammer
\langle proof \rangle
lemma True nitpick[satisfy] \langle proof \rangle
     Modal collapse and monotheism: not implied.
lemma MC: |\forall \Phi.(\Phi \rightarrow \Box \Phi)| \text{ nitpick } \langle proof \rangle
lemma MT: [\forall x \ y.(((\mathcal{G} \ x) \land (\mathcal{G} \ y)) \rightarrow (x=y))]
  nitpick \langle proof \rangle
     Gödel's A1, A4, A5: not implied anymore.
lemma A1: |\forall X.((\neg(\mathcal{P} X)) \leftrightarrow (\mathcal{P}(\rightarrow X)))| \text{ nitpick } \langle proof \rangle
lemma A_4: |\forall X.((\mathcal{P} X) \to \Box(\mathcal{P} X))| \text{ nitpick } \langle proof \rangle
lemma A5: |\mathcal{P}| \mathcal{NE}| nitpick \langle proof \rangle
     Checking filter and ultrafilter properties.
theorem F1: |Filter \mathcal{P}| \langle proof \rangle
theorem U1: \lfloor UFilter \mathcal{P} \rfloor nitpick \langle proof \rangle
end
          Simplified Variant with Axiom T2 (Fig. 7 in [4])
4.6
theory SimpleVariantPG imports
  HOML
  MFilter
  BaseDefs
begin
     Axiom's of simplified variant with A3 replaced.
axiomatization where
  A1': |\neg(\mathcal{P}(\lambda x.(x\neq x)))| and
  A2': |\forall X Y.(((\mathcal{P} X) \land ((X \sqsubseteq Y) \lor (X \Rrightarrow Y))) \rightarrow (\mathcal{P} Y))| and
```

```
T2: [\mathcal{P} \mathcal{G}]
     Necessary existence of a Godlike entity.
theorem T6: |\Box(\exists^E \mathcal{G})| — Proof found by sledgehammer
\langle proof \rangle
lemma True nitpick[satisfy] \langle proof \rangle
     Modal collapse and Monotheism: not implied.
lemma MC: |\forall \Phi.(\Phi \rightarrow \Box \Phi)| \text{ nitpick } \langle proof \rangle
lemma MT: [\forall x \ y.(((\mathcal{G} \ x) \land (\mathcal{G} \ y)) \rightarrow (x=y))] nitpick \langle proof \rangle
end
4.7
          Simplified Variant with Simple Entailment in Logic K
          (Fig. 8 in [4])
theory SimpleVariantSE imports
  HOML
  MFilter
  BaseDefs
begin
     Axiom's of new variant based on ultrafilters.
axiomatization where
  A1': |\neg(\mathcal{P}(\lambda x.(x\neq x)))| and
  A2'': |\forall X Y.(((\mathcal{P} X) \land (X \sqsubseteq Y)) \rightarrow (\mathcal{P} Y))| and
  T2: |\mathcal{P} \mathcal{G}|
     Necessary existence of a Godlike entity.
theorem T6: \lfloor \Box (\exists^E \mathcal{G}) \rfloor \langle proof \rangle
theorem T7: \lfloor \exists^E \mathcal{G} \rfloor \langle proof \rangle
     Possible existence of a Godlike: has counterodel.
lemma T3: |\diamondsuit(\exists E \mathcal{G})| \text{ nitpick } \langle proof \rangle
lemma T3': assumes T: |\forall \varphi.((\Box \varphi) \rightarrow \varphi)|
  shows [\lozenge(\exists^E \mathcal{G})]
  \langle proof \rangle
\quad \mathbf{end} \quad
          Simplified Variant with Simple Entailment in Logic T
4.8
          (Fig. 9 in [4])
theory Simple Variant SEin T imports
```

Axiom's of new variant based on ultrafilters.

HOML MFilter BaseDefs begin

```
axiomatization where
  A1': |\neg(\mathcal{P}(\lambda x.(x\neq x)))| and
  A2'': |\forall X Y.(((\mathcal{P} X) \land (X \sqsubseteq Y)) \rightarrow (\mathcal{P} Y))| and
  T2: |\mathcal{P} \mathcal{G}|
     Modal Logic T.
axiomatization where T: |\forall \varphi.((\Box \varphi) \rightarrow \varphi)|
lemma T': |\forall \varphi.(\varphi \to (\Diamond \varphi))| \langle proof \rangle
     Necessary existence of a Godlike entity.
theorem T6: |\Box(\exists^E \mathcal{G})| — Proof found by sledgehammer
\langle proof \rangle
     T6 again, with an alternative, simpler proof.
theorem T6again: |\Box(\exists^E \mathcal{G})|
\langle proof \rangle
end
         Hauptfiltervariant (Fig. 10 in [4])
4.9
theory SimpleVariantHF imports
  HOML
  MFilter
  BaseDefs
begin
     Definition: Set of supersets of X, we call this \mathcal{HF} X.
abbreviation HF::\gamma \Rightarrow (\gamma \Rightarrow \sigma) where HF X \equiv \lambda Y.(X \sqsubseteq Y)
     Postulate: \mathcal{HF} \mathcal{G} is a filter; i.e., Hauptfilter of \mathcal{G}.
axiomatization where F1: |Filter(HF \mathcal{G})|
     Necessary existence of a Godlike entity.
theorem T6: \lfloor \Box(\exists^E \mathcal{G}) \rfloor \langle proof \rangle
theorem T6again: |\Box(\exists E \mathcal{G})|
\langle proof \rangle
     Possible existence of Godlike entity not implied.
lemma T3: |\diamondsuit(\exists^E \mathcal{G})| \text{ nitpick } \langle proof \rangle
     Axiom T enforces possible existence of Godlike entity.
axiomatization
lemma T3: assumes T: [\forall \varphi.((\Box \varphi) \rightarrow \varphi)]
           shows |\diamondsuit(\exists^E \mathcal{G})| \langle proof \rangle
lemma True nitpick[satisfy] \langle proof \rangle
```

Modal collapse: not implied anymore.

```
lemma MT: [\forall x \ y.(((\mathcal{G} \ x) \land (\mathcal{G} \ y)) \rightarrow (x=y))]
             nitpick \langle proof \rangle
end
             Formal Study of Version No.2 of Gödel's Argument as
4.10
              Reported by Kanckos and Lethen, 2019 [6] (Fig. 11 in
              [4])
theory KanckosLethenNo2Possibilist imports
   HOML
   MFilter
   BaseDefs
begin
      Axioms of Version No. 2 [6].
abbreviation delta (\langle \Delta \rangle) where \Delta A \equiv \lambda x.(\forall \psi. ((A \psi) \rightarrow (\psi x)))
abbreviation N(\langle \mathcal{N} \rangle) where \mathcal{N} \varphi \equiv \lambda x.(\Box(\varphi x))
axiomatization where
   Axiom1: |\forall \varphi \psi.(((\mathcal{P} \varphi) \land (\Box(\forall x. ((\varphi x) \rightarrow (\psi x))))) \rightarrow (\mathcal{P} \psi))| and — The
\square can be omitted here; the proofs still work.
  Axiom2: |\forall A . (\Box((\forall \varphi.((A \varphi) \to (\mathcal{P} \varphi))) \to (\mathcal{P} (\Delta A))))| and — The \Box can
be omitted here; the proofs still work.
   Axiom3: |\forall \varphi.((\mathcal{P} \varphi) \to (\mathcal{P} (\mathcal{N} \varphi)))| and
  Axiom4: |\forall \varphi.((\mathcal{P} \varphi) \to (\neg(\mathcal{P}(\neg\varphi))))| and
    Logic S5
  axB: |\forall \varphi.(\varphi \to \Box \Diamond \varphi)| \text{ and } axM: |\forall \varphi.((\Box \varphi) \to \varphi)| \text{ and } ax4: |\forall \varphi.((\Box \varphi) \to \varphi)|
(\Box\Box\varphi))
      Sahlqvist correspondences: they are better suited for proof automation.
lemma axB': \forall x \ y. \ \neg(x\mathbf{r}y) \lor (y\mathbf{r}x) \ \langle proof \rangle
lemma axM': \forall x. (xrx) \langle proof \rangle
lemma ax4': \forall x \ y \ z. (((xry) \land (yrz)) \longrightarrow (xrz)) \langle proof \rangle
      Proofs for all theorems for No.2 from [6].
theorem Theorem 0: |\forall \varphi \ \psi.((\forall Q.\ ((Q \ \varphi) \ \rightarrow (Q \ \psi)))) \rightarrow \ ((\mathcal{P} \ \varphi) \rightarrow (\mathcal{P} \ \varphi)))|
\langle proof \rangle
theorem 1: |P G| \langle proof \rangle
theorem Theorem2: |\forall x. ((\mathcal{G} x) \rightarrow (\exists y. \mathcal{G} y))| \langle proof \rangle
theorem Theorem3a: |\mathcal{P}(\lambda x. (\exists y. \mathcal{G}y))| \langle proof \rangle
theorem Theorem3b: [\Box(\mathcal{P}(\lambda x.(\Box(\exists y. \mathcal{G} y)))))] \langle proof \rangle
theorem Theorem 4: |\forall x. \Box((\mathcal{G} x) \rightarrow ((\mathcal{P} (\lambda x. (\Box(\exists y. \mathcal{G} y))))) \rightarrow (\Box(\exists y. \mathcal{G} y))))|
\langle proof \rangle
theorem Theorem 5: [\forall x. \ \Box((\mathcal{G} \ x) \rightarrow (\Box(\exists y. \ \mathcal{G} \ y)))] \ \langle proof \rangle
theorem Theorem6: [\Box((\exists y. \mathcal{G} y) \rightarrow (\Box(\exists y. \mathcal{G} y)))] \langle proof \rangle
theorem Theorem 7: |\Box((\Diamond(\exists y. \mathcal{G} y)) \rightarrow (\Box(\exists y. \mathcal{G} y)))| \langle proof \rangle
theorem Theorem 8: |\Box(\exists y. \mathcal{G} y)| \langle proof \rangle
theorem Theorem 9: |\forall \varphi. ((\mathcal{P} \varphi) \rightarrow \Diamond(\exists x. \varphi x))| \langle proof \rangle
```

lemma $MC: |\forall \Phi.(\Phi \rightarrow \Box \Phi)| \text{ nitpick } \langle proof \rangle$

Short proof of Theorem8; analogous to the one presented in Sec. 7 of Benzmüller 2020.

theorem $[\Box(\exists y. \mathcal{G} y)]$ — Note: this version of the proof uses only axB' and axM'. $\langle proof \rangle$

```
theorem T5: \lfloor (\diamondsuit(\exists y. \mathcal{G} y)) \rightarrow \Box(\exists y. \mathcal{G} y) \rfloor — Obvious: If we can prove Theorem8, then we also have T5. \langle proof \rangle
```

Another short proof of Theorem8.

```
theorem [\Box(\exists y. \mathcal{G} y)] — Note: fewer assumptions used in some cases than in [6]. \langle proof \rangle
```

Are the axioms of the simplified versions implied?

Actualist version of the axioms.

```
lemma A1': \lfloor \neg (\mathcal{P}(\lambda x.(x \neq x))) \rfloor \langle proof \rangle lemma A2': \lfloor \forall X \ Y.(((\mathcal{P} \ X) \land ((X \sqsubseteq Y) \lor (X \Rrightarrow Y))) \rightarrow (\mathcal{P} \ Y)) \rfloor nitpick \langle proof \rangle lemma A3: \lfloor \forall \mathcal{Z}.((\mathcal{P} \text{os } \mathcal{Z}) \rightarrow (\forall X.((X \sqcap \mathcal{Z}) \rightarrow (\mathcal{P} \ X)))) \rfloor nitpick \langle proof \rangle
```

Possibilist version of the axioms.

```
abbreviation a\ (\cdot - \sqsubseteq^p - \cdot) where X \sqsubseteq^p Y \equiv \forall z.((X\ z) \to (Y\ z)) abbreviation b\ (\cdot - \Rrightarrow^p - \cdot) where X \Rrightarrow^p Y \equiv \Box(X \sqsubseteq^p Y) abbreviation d\ (\cdot - \Box^p - \cdot) where X \Box^p \mathcal{Z} \equiv \Box(\forall\ u.((X\ u) \leftrightarrow (\forall\ Y.((\mathcal{Z}\ Y) \to (Y\ u)))))
```

```
\begin{array}{l} \mathbf{lemma} \ A1'P \colon \left \lfloor \neg (\mathcal{P}(\lambda x.(x \neq x))) \right \rfloor \ \langle proof \rangle \\ \mathbf{lemma} \ A2'P \colon \left \lfloor \forall \ X \ Y.(((\mathcal{P} \ X) \ \land \ ((X \sqsubseteq^p Y) \lor (X \Rrightarrow^p Y))) \ \rightarrow \ (\mathcal{P} \ Y)) \right \rfloor \ \langle proof \rangle \\ \mathbf{lemma} \ A2'aP \colon \left \lfloor \forall \ X \ Y.(((\mathcal{P} \ X) \ \land \ (X \Rrightarrow^p Y)) \ \rightarrow \ (\mathcal{P} \ Y)) \right \rfloor \ \langle proof \rangle \\ \mathbf{lemma} \ A2'bP \colon \left \lfloor \forall \ X \ Y.(((\mathcal{P} \ X) \ \land \ (X \sqsubseteq^p Y)) \ \rightarrow \ (\mathcal{P} \ Y)) \right \rfloor \ \langle proof \rangle \\ \mathbf{lemma} \ A3P \colon \left \lfloor \forall \ \mathcal{Z}.((\mathcal{P} \text{os } \mathcal{Z}) \ \rightarrow \ (\forall \ X.((X \sqcap^p \mathcal{Z}) \ \rightarrow \ (\mathcal{P} \ X))))) \right \rfloor \\ \langle proof \rangle \end{array}
```

Are Axiom2 and A3 equivalent? Only when assuming Axiom1 and axiom \mathcal{M} .

```
\begin{array}{ll} \mathbf{lemma} & \left \lfloor \forall \ A \ . (\Box((\forall \ \varphi.((A \ \varphi) \ \rightarrow \ (\mathcal{P} \ \varphi))) \ \rightarrow \ (\mathcal{P} \ (\Delta \ A)))) \right \rfloor \ \equiv \ \left \lfloor \forall \ \mathcal{Z}.((\mathcal{P} \mathrm{os} \ \mathcal{Z}) \ \rightarrow \ (\forall \ X.((X \ \sqcap^{p}\mathcal{Z}) \ \rightarrow \ (\mathcal{P} \ X)))) \right \rfloor \ \\ & \langle \mathit{proof} \, \rangle \\ & \mathbf{end} \end{array}
```

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