

Exploring Simplified Variants of Gödel's Ontological Argument in Isabelle/HOL

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December 14, 2021

Abstract

Simplified variants of Gödel's ontological argument are explored. Among those is a particularly interesting simplified argument which is (i) valid already in basic modal logics K or KT, (ii) which does not suffer from modal collapse, and (iii) which avoids the rather complex predicates of essence (Ess.) and necessary existence (NE) as used by Gödel.

Whether the presented variants increase or decrease the attractiveness and persuasiveness of the ontological argument is a question I would like to pass on to philosophy and theology.

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1 Background Reading

The selected simplified variants of Gödel's ontological argument [5, 7] as presented in §3 have first been extracted from the insights gained in §6 of [4]. These variants are also influenced by the work presented in [1] and they significantly extend the findings from [3]. In §4 we additionally include the sources from [4].

2 Higher-Order Modal Logic in HOL (cf. [2] and Fig. 1 in [4]).

```
theory HOML imports Main
begin
nitpick-params[user-axioms,expect=genuine]
```

Type *i* is associated with possible worlds and type *e* with entities:

```
typedecl i — Possible worlds
typedecl e — Individuals
type-synonym  $\sigma = i \Rightarrow bool$  — World-lifted propositions
type-synonym  $\gamma = e \Rightarrow \sigma$  — Lifted predicates
type-synonym  $\mu = \sigma \Rightarrow \sigma$  — Unary modal connectives
type-synonym  $\nu = \sigma \Rightarrow \sigma \Rightarrow \sigma$  — Binary modal connectives
```

Logical connectives (operating on truth-sets):

```
abbreviation c1:: $\sigma$  ( $\perp$ ) where  $\perp \equiv \lambda w. False$ 
abbreviation c2:: $\sigma$  ( $\top$ ) where  $\top \equiv \lambda w. True$ 
abbreviation c3:: $\mu$  ( $\neg$ [52]53) where  $\neg\varphi \equiv \lambda w. \neg(\varphi w)$ 
abbreviation c4:: $\nu$  (infix $\wedge$ 50) where  $\varphi \wedge \psi \equiv \lambda w. (\varphi w) \wedge (\psi w)$ 
abbreviation c5:: $\nu$  (infix $\vee$ 49) where  $\varphi \vee \psi \equiv \lambda w. (\varphi w) \vee (\psi w)$ 
abbreviation c6:: $\nu$  (infix $\rightarrow$ 48) where  $\varphi \rightarrow \psi \equiv \lambda w. (\varphi w) \rightarrow (\psi w)$ 
abbreviation c7:: $\nu$  (infix $\leftrightarrow$ 47) where  $\varphi \leftrightarrow \psi \equiv \lambda w. (\varphi w) \leftrightarrow (\psi w)$ 
consts R:: $i \Rightarrow i \Rightarrow bool$  (-r-) — Accessibility relation
abbreviation c8:: $\mu$  ( $\Box$ [54]55) where  $\Box\varphi \equiv \lambda w. \forall v. (wr v) \rightarrow (\varphi v)$ 
abbreviation c9:: $\mu$  ( $\Diamond$ [54]55) where  $\Diamond\varphi \equiv \lambda w. \exists v. (wr v) \wedge (\varphi v)$ 
abbreviation c10:: $\gamma \Rightarrow \gamma$  ( $\neg$ [52]53) where  $\neg\Phi \equiv \lambda x. \lambda w. \neg(\Phi x w)$ 
abbreviation c11:: $\gamma \Rightarrow \gamma$  ( $\rightarrow$ -) where  $\rightarrow\Phi \equiv \lambda x. \lambda w. \neg(\Phi x w)$ 
abbreviation c12:: $e \Rightarrow e \Rightarrow \sigma$  ( $=$ -) where  $x=y \equiv \lambda w. (x=y)$ 
abbreviation c13:: $e \Rightarrow e \Rightarrow \sigma$  ( $\neq$ -) where  $x \neq y \equiv \lambda w. (x \neq y)$ 
```

Polymorphic possibilist quantification:

abbreviation $q1::('a\Rightarrow\sigma)\Rightarrow\sigma$ (\forall) **where** $\forall\Phi\equiv\lambda w.\forall x.(\Phi x w)$
abbreviation $q2$ (**binder** \forall [10]11) **where** $\forall x.\varphi(x)\equiv\forall\varphi$
abbreviation $q3::('a\Rightarrow\sigma)\Rightarrow\sigma$ (\exists) **where** $\exists\Phi\equiv\lambda w.\exists x.(\Phi x w)$
abbreviation $q4$ (**binder** \exists [10]11) **where** $\exists x.\varphi(x)\equiv\exists\varphi$

Actualist quantification for individuals/entities:

consts $existsAt::\gamma$ ($-\textcircled{-}$)
abbreviation $q5::\gamma\Rightarrow\sigma$ (\forall^E) **where** $\forall^E\Phi\equiv\lambda w.\forall x.(x\textcircled{w})\longrightarrow(\Phi x w)$
abbreviation $q6$ (**binder** \forall^E [8]9) **where** $\forall^E x.\varphi(x)\equiv\forall^E\varphi$
abbreviation $q7::\gamma\Rightarrow\sigma$ (\exists^E) **where** $\exists^E\Phi\equiv\lambda w.\exists x.(x\textcircled{w})\wedge(\Phi x w)$
abbreviation $q8$ (**binder** \exists^E [8]9) **where** $\exists^E x.\varphi(x)\equiv\exists^E\varphi$

Meta-logical predicate for global validity:

abbreviation $g1::\sigma\Rightarrow bool$ ($[-]$) **where** $[\psi]\equiv\forall w.\psi w$

Barcan and converse Barcan formula:

lemma *True nitpick[satisfy] oops* — Model found by Nitpick
lemma $[(\forall^E x.\Box(\varphi x))\rightarrow\Box(\forall^E x.(\varphi x))]$ **nitpick oops** — Ctm
lemma $[\Box(\forall^E x.(\varphi x))\rightarrow(\forall^E x.\Box(\varphi x))]$ **nitpick oops** — Ctm
lemma $[(\forall x.\Box(\varphi x))\rightarrow\Box(\forall x.\varphi x)]$ **by simp**
lemma $[\Box(\forall x.(\varphi x))\rightarrow(\forall x.\Box(\varphi x))]$ **by simp**
end
theory *DisableKodkodScala*
imports *Main*
begin

Some of the nitpick invocation within this AFP entry do not work, if "Kodkod Scala" is enabled, i.e., if the box under Plugin Options — Isabelle — General — Miscellaneous Tools — Kodkod Scala is activated. Therefore, in this theory we explicitly disable this configuration option.

ML \langle
Options.default-put-bool **system-option** $\langle kodkod-scala \rangle$ *false*
 \rangle

end

3 Selected Simplified Ontological Argument

theory *SimplifiedOntologicalArgument* **imports**
HOML
DisableKodkodScala
begin

Positive properties:

consts $posProp::\gamma\Rightarrow\sigma$ (\mathcal{P})

An entity x is God-like if it possesses all positive properties.

definition G (\mathcal{G}) **where** $\mathcal{G}(x)\equiv\forall\Phi.(\mathcal{P}(\Phi)\rightarrow\Phi(x))$

The axiom's of the simplified variant are presented next; these axioms are further motivated in [4, 1]).

Self-difference is not a positive property (possible alternative: the empty property is not a positive property).

axiomatization where CORO1: $[\neg(\mathcal{P}(\lambda x.(x \neq x)))]$

A property entailed by a positive property is positive.

axiomatization where CORO2: $[\forall \Phi \Psi. \mathcal{P}(\Phi) \wedge (\forall x. \Phi(x) \rightarrow \Psi(x)) \rightarrow \mathcal{P}(\Psi)]$

Being Godlike is a positive property.

axiomatization where AXIOM3: $[\mathcal{P} \mathcal{G}]$

3.1 Verifying the Selected Simplified Ontological Argument (version 1)

The existence of a non-exemplified positive property implies that self-difference (or, alternatively, the empty property) is a positive property.

lemma LEMMA1: $[(\exists \Phi. (\mathcal{P}(\Phi) \wedge \neg(\exists x. \Phi(x)))) \rightarrow \mathcal{P}(\lambda x.(x \neq x))]$
using CORO2 by meson

A non-exemplified positive property does not exist.

lemma LEMMA2: $[\neg(\exists \Phi. (\mathcal{P}(\Phi) \wedge \neg(\exists x. \Phi(x))))]$
using CORO1 LEMMA1 by blast

Positive properties are exemplified.

lemma LEMMA3: $[\forall \Phi. (\mathcal{P}(\Phi) \rightarrow (\exists x. \Phi(x)))]$
using LEMMA2 by blast

There exists a God-like entity.

theorem THEOREM3': $[\exists x. \mathcal{G}(x)]$
using AXIOM3 LEMMA3 by auto

Necessarily, there exists a God-like entity.

theorem THEOREM3: $[\Box(\exists x. \mathcal{G}(x))]$
using THEOREM3' by simp

However, the possible existence of Godlike entity is not implied.

theorem CORO: $[\Diamond(\exists x. \mathcal{G}(x))]$
nitpick oops

3.2 Verifying the Selected Simplified Ontological Argument (version 2)

We switch to logic T.

axiomatization where T: $[\forall \varphi. \Box \varphi \rightarrow \varphi]$

lemma T': $[\forall \varphi. \varphi \rightarrow \Diamond \varphi]$ **using T bymetis**

Positive properties are possibly exemplified.

theorem *THEOREM1*: $[\forall \Phi. \mathcal{P}(\Phi) \rightarrow \diamond(\exists x. \Phi(x))]$
using *CORO1 CORO2 T'* by *metis*

Possibly there exists a God-like entity.

theorem *CORO*: $[\diamond(\exists x. \mathcal{G}(x))]$
using *AXIOM3 THEOREM1* by *auto*

The possible existence of a God-like entity implies the necessary existence of a God-like entity.

theorem *THEOREM2*: $[\diamond(\exists x. \mathcal{G}(x)) \rightarrow \Box(\exists x. \mathcal{G}(x))]$
using *AXIOM3 CORO1 CORO2* by *metis*

Necessarily, there exists a God-like entity.

theorem *THEO3*: $[\Box(\exists x. \mathcal{G}(x))]$
using *CORO THEOREM2* by *blast*

There exists a God-like entity.

theorem *THEO3'*: $[\exists x. \mathcal{G}(x)]$
using *T THEO3* by *metis*

Modal collapse is not implied; nitpick reports a countermodel.

lemma *MC*: $[\forall \Phi. \Phi \rightarrow \Box \Phi]$ **nitpick oops**

Consistency of the theory; nitpick reports a model.

lemma *True nitpick[satisfy] oops*
end

4 Presentation of All Variants as Studied in [4]

4.1 Preliminaries: Modal Ultrafilter (Fig. 2 in [4])

theory *MFilter* **imports** *HOML*
begin

Some abbreviations for auxiliary operations.

abbreviation *a*:: $\gamma \Rightarrow (\gamma \Rightarrow \sigma) \Rightarrow \sigma$ ($-\in-$) **where** $x \in S \equiv S x$

abbreviation *b*:: γ (\emptyset) **where** $\emptyset \equiv \lambda x. \perp$

abbreviation *c*:: γ (\mathbf{U}) **where** $\mathbf{U} \equiv \lambda x. \top$

abbreviation *d*:: $\gamma \Rightarrow \gamma \Rightarrow \sigma$ ($-\subseteq-$) **where** $\varphi \subseteq \psi \equiv \forall x. ((\varphi x) \rightarrow (\psi x))$

abbreviation *e*:: $\gamma \Rightarrow \gamma \Rightarrow \gamma$ ($-\sqcap-$) **where** $\varphi \sqcap \psi \equiv \lambda x. ((\varphi x) \wedge (\psi x))$

abbreviation *f*:: $\gamma \Rightarrow \gamma$ ($^{-1}-$) **where** $^{-1}\psi \equiv \lambda x. \neg(\psi x)$

Definition of modal filter.

abbreviation *g*:: $(\gamma \Rightarrow \sigma) \Rightarrow \sigma$ (*Filter*)

where *Filter* $\Phi \equiv (((\mathbf{U} \in \Phi) \wedge \neg(\emptyset \in \Phi))$

$\wedge (\forall \varphi \psi. (((\varphi \in \Phi) \wedge (\varphi \subseteq \psi)) \rightarrow (\psi \in \Phi))))$

$\wedge (\forall \varphi \psi. (((\varphi \in \Phi) \wedge (\psi \in \Phi)) \rightarrow ((\varphi \sqcap \psi) \in \Phi)))$

Definition of modal ultrafilter .

abbreviation $h::(\gamma \Rightarrow \sigma) \Rightarrow \sigma$ (*UFilter*) **where**
 $UFilter \Phi \equiv (Filter \Phi) \wedge (\forall \varphi. ((\varphi \in \Phi) \vee ((^{-1}\varphi) \in \Phi)))$

Modal filter and modal ultrafilter are consistent.

lemma $[\forall \Phi \varphi. ((UFilter \Phi) \rightarrow \neg((\varphi \in \Phi) \wedge ((^{-1}\varphi) \in \Phi)))]$ **by force**
end

4.2 Preliminaries: Further Basic Notions (Fig. 3 in [4])

theory *BaseDefs* **imports** *HOML*
begin

Positive properties.

consts $posProp::\gamma \Rightarrow \sigma$ (\mathcal{P})

Basic definitions for modal ontological argument.

abbreviation a ($-\sqsubseteq$) **where** $X \sqsubseteq Y \equiv \forall z. ((X z) \rightarrow (Y z))$

abbreviation b ($-\Rightarrow$) **where** $X \Rightarrow Y \equiv \Box(X \sqsubseteq Y)$

abbreviation c ($\mathcal{P}os$) **where** $\mathcal{P}os Z \equiv \forall X. ((Z X) \rightarrow (\mathcal{P} X))$

abbreviation d ($-\sqcap$) **where** $X \sqcap Z \equiv \Box(\forall u. ((X u) \leftrightarrow (\forall Y. ((Z Y) \rightarrow (Y u))))))$

Definition of Godlike.

definition G (\mathcal{G}) **where** $\mathcal{G} x \equiv \forall Y. ((\mathcal{P} Y) \rightarrow (Y x))$

Definitions of Essence and Necessary Existence.

definition E (\mathcal{E}) **where** $\mathcal{E} Y x \equiv (Y x) \wedge (\forall Z. ((Z x) \rightarrow (Y \Rightarrow Z)))$

definition NE ($\mathcal{N}\mathcal{E}$) **where** $\mathcal{N}\mathcal{E} x \equiv \forall Y. ((\mathcal{E} Y x) \rightarrow \Box(\exists^E Y))$

end

4.3 Ultrafilter Analysis of Scott's Variant (Fig. 3 in [4])

theory *ScottVariant* **imports**

HOML

MFilter

BaseDefs

begin

Axioms of Scott's variant.

axiomatization where

$A1: [\forall X. ((\neg(\mathcal{P} X)) \leftrightarrow (\mathcal{P}(\neg X)))]$ **and**

$A2: [\forall X Y. (((\mathcal{P} X) \wedge (X \Rightarrow Y)) \rightarrow (\mathcal{P} Y))]$ **and**

$A3: [\forall Z. ((\mathcal{P}os Z) \rightarrow (\forall X. ((X \sqcap Z) \rightarrow (\mathcal{P} X))))]$ **and**

$A4: [\forall X. ((\mathcal{P} X) \rightarrow \Box(\mathcal{P} X))]$ **and**

$A5: [\mathcal{P} \mathcal{N}\mathcal{E}]$ **and**

$B: [\forall \varphi. (\varphi \rightarrow \Box \Diamond \varphi)]$ — Logic KB

lemma B' : $\forall x y. \neg(x y) \vee (y x)$ **using** B **by fastforce**

Necessary existence of a Godlike entity.

theorem T6: $[\Box(\exists^E \mathcal{G})]$

proof –

have T1: $[\forall X.((\mathcal{P} X) \rightarrow \Diamond(\exists^E X))]$ **using A1 A2 by blast**

have T2: $[\mathcal{P} \mathcal{G}]$ **by (metis A3 G-def)**

have T3: $[\Diamond(\exists^E \mathcal{G})]$ **using T1 T2 by simp**

have T4: $[\forall^E x.((\mathcal{G} x) \rightarrow (\mathcal{E} \mathcal{G} x))]$ **unfolding G-def E-def using A1 A4 by metis**

have T5: $[(\Diamond(\exists^E \mathcal{G})) \rightarrow \Box(\exists^E \mathcal{G})]$ **by (smt A5 G-def B' NE-def T4)**

thus ?thesis using T3 by blast qed

Existence of a Godlike entity.

lemma $[\exists^E \mathcal{G}]$ **using A1 A2 B' T6 by blast**

Consistency

lemma True nitpick[satisfy] oops — Model found.

Modal collapse: holds.

lemma MC: $[\forall \Phi.(\Phi \rightarrow \Box \Phi)]$

proof – {

fix w fix Q

have 1: $[\forall x.((\mathcal{G} x w) \rightarrow (\forall Z.((Z x) \rightarrow \Box(\forall^E z.((\mathcal{G} z) \rightarrow (Z z)))))) w]$
by (metis A1 A4 G-def)

have 2: $[(\exists x. \mathcal{G} x w) \rightarrow ((Q \rightarrow \Box(\forall^E z.((\mathcal{G} z) \rightarrow Q))) w)]$
using 1 by force

have 3: $(Q \rightarrow \Box Q) w$ **using B' T6 2 by blast}**

thus ?thesis by auto qed

Analysis of positive properties using ultrafilters.

theorem U1: $[UFilter \mathcal{P}]$ — Proof found by sledgehammer

proof –

have 1: $[(\bigcup \in \mathcal{P}) \wedge \neg(\emptyset \in \mathcal{P})]$ **using A1 A2 by blast**

have 2: $[\forall \varphi \psi.(((\varphi \in \mathcal{P}) \wedge (\varphi \subseteq \psi)) \rightarrow (\psi \in \mathcal{P}))]$ **by (smt A2 B' MC)**

have 3: $[\forall \varphi \psi.(((\varphi \in \mathcal{P}) \wedge (\psi \in \mathcal{P})) \rightarrow ((\varphi \cap \psi) \in \mathcal{P}))]$ **by (metis A1 A2 G-def B' T6)**

have 4: $[\forall \varphi.((\varphi \in \mathcal{P}) \vee ((^{-1}\varphi) \in \mathcal{P}))]$ **using A1 by blast**

thus ?thesis using 1 2 3 4 by simp qed

lemma L1: $[\forall X Y.((X \Rightarrow Y) \rightarrow (X \sqsubseteq Y))]$ **by (metis A1 A2 MC)**

lemma L2: $[\forall X Y.(((\mathcal{P} X) \wedge (X \sqsubseteq Y)) \rightarrow (\mathcal{P} Y))]$ **by (smt A2 B' MC)**

Set of supersets of X, we call this HF X.

abbreviation HF **where** $HF X \equiv \lambda Y.(X \sqsubseteq Y)$

$HF \mathcal{G}$ is a filter; hence, $HF \mathcal{G}$ is Hauptfilter of \mathcal{G} .

lemma F1: $[Filter (HF \mathcal{G})]$ **by (metis A2 B' T6 U1)**

lemma F2: $[UFilter (HF \mathcal{G})]$ **by (smt A1 F1 G-def)**

T6 follows directly from F1.

theorem T6again: $[\Box(\exists^E \mathcal{G})]$ **using F1 by simp**

end

4.4 Ultrafilter Variant (Fig. 5 in [4])

theory *UFilterVariant* **imports**

HOML

MFilter

BaseDefs

DisableKodkodScala

begin

Axiom's of ultrafilter variant.

axiomatization where

U1: [*UFilter* \mathcal{P}] **and**

A2: [$\forall X Y. ((\mathcal{P} X) \wedge (X \Rightarrow Y)) \rightarrow (\mathcal{P} Y)$] **and**

A3: [$\forall \mathcal{Z}. ((\mathcal{P} \text{os } \mathcal{Z}) \rightarrow (\forall X. ((X \sqcap \mathcal{Z}) \rightarrow (\mathcal{P} X))))$]

Necessary existence of a Godlike entity.

theorem *T6*: [$\Box(\exists^E \mathcal{G})$] — Proof also found by sledgehammer

proof –

have *T1*: [$\forall X. ((\mathcal{P} X) \rightarrow \Diamond(\exists^E X))$] **by** (*metis A2 U1*)

have *T2*: [$\mathcal{P} \mathcal{G}$] **by** (*metis A3 G-def*)

have *T3*: [$\Diamond(\exists^E \mathcal{G})$] **using** *T1 T2* **by** *simp*

have *T5*: [$(\Diamond(\exists^E \mathcal{G})) \rightarrow \Box(\exists^E \mathcal{G})$] **by** (*metis A2 G-def T2 U1*)

thus *?thesis* **using** *T3* **by** *blast qed*

Checking for consistency.

lemma *True nitpick[satisfy] oops* — Model found

Checking for modal collapse.

lemma *MC*: [$\forall \Phi. (\Phi \rightarrow \Box \Phi)$] **nitpick oops** — Countermodel
end

4.5 Simplified Variant (Fig. 6 in [4])

theory *SimpleVariant* **imports**

HOML

MFilter

BaseDefs

DisableKodkodScala

begin

Axiom's of new, simplified variant.

axiomatization where

A1': [$\neg(\mathcal{P}(\lambda x. (x \neq x)))$] **and**

A2': [$\forall X Y. ((\mathcal{P} X) \wedge ((X \sqsubseteq Y) \vee (X \Rightarrow Y))) \rightarrow (\mathcal{P} Y)$] **and**

A3: [$\forall \mathcal{Z}. ((\mathcal{P} \text{os } \mathcal{Z}) \rightarrow (\forall X. ((X \sqcap \mathcal{Z}) \rightarrow (\mathcal{P} X))))$]

lemma *T2*: $[\mathcal{P} \mathcal{G}]$ **by** (*metis A3 G-def*) — From A3
lemma *L1*: $[\mathcal{P}(\lambda x.(x=x))]$ **by** (*metis A2' A3*)

Necessary existence of a Godlike entity.

theorem *T6*: $[\Box(\exists^E \mathcal{G})]$ — Proof found by sledgehammer
proof –

have *T1*: $[\forall X.((\mathcal{P} X) \rightarrow \Diamond(\exists^E X))]$ **by** (*metis A1' A2'*)
have *T3*: $[\Diamond(\exists^E \mathcal{G})]$ **using** *T1 T2* **by** *simp*
have *T5*: $[(\Diamond(\exists^E \mathcal{G})) \rightarrow \Box(\exists^E \mathcal{G})]$ **by** (*metis A1' A2' T2*)
thus *?thesis* **using** *T3* **by** *blast qed*

lemma *True nitpick* $[satisfy]$ **oops** — Consistency: model found

Modal collapse and monotheism: not implied.

lemma *MC*: $[\forall \Phi.(\Phi \rightarrow \Box \Phi)]$ **nitpick oops** — Countermodel

lemma *MT*: $[\forall x y.(((\mathcal{G} x) \wedge (\mathcal{G} y)) \rightarrow (x=y))]$

nitpick oops — Countermodel.

Gödel's A1, A4, A5: not implied anymore.

lemma *A1*: $[\forall X.((\neg(\mathcal{P} X)) \leftrightarrow (\mathcal{P}(\neg X)))]$ **nitpick oops** — Countermodel

lemma *A4*: $[\forall X.((\mathcal{P} X) \rightarrow \Box(\mathcal{P} X))]$ **nitpick oops** — Countermodel

lemma *A5*: $[\mathcal{P} \mathcal{N}\mathcal{E}]$ **nitpick oops** — Countermodel

Checking filter and ultrafilter properties.

theorem *F1*: $[Filter \mathcal{P}]$ **oops** — Proof found by sledgehammer, but reconstruction timeout

theorem *U1*: $[UFilter \mathcal{P}]$ **nitpick oops** — Countermodel

end

4.6 Simplified Variant with Axiom T2 (Fig. 7 in [4])

theory *SimpleVariantPG* **imports**

HOML

MFilter

BaseDefs

DisableKodkodScala

begin

Axiom's of simplified variant with A3 replaced.

axiomatization where

A1': $[\neg(\mathcal{P}(\lambda x.(x \neq x)))]$ **and**

A2': $[\forall X Y.(((\mathcal{P} X) \wedge ((X \sqsubseteq Y) \vee (X \ni Y))) \rightarrow (\mathcal{P} Y))]$ **and**

T2: $[\mathcal{P} \mathcal{G}]$

Necessary existence of a Godlike entity.

theorem *T6*: $[\Box(\exists^E \mathcal{G})]$ — Proof found by sledgehammer
proof –

have *T1*: $[\forall X.((\mathcal{P} X) \rightarrow \Diamond(\exists^E X))]$ **by** (*metis A1' A2'*)

have *T3*: $[\Diamond(\exists^E \mathcal{G})]$ **using** *T1 T2* **by** *simp*

have $T5$: $[(\diamond(\exists^E \mathcal{G})) \rightarrow \Box(\exists^E \mathcal{G})]$ **by** (*metis* $A1'$ $A2'$ $T2$)
thus *?thesis* **using** $T3$ **by** *blast qed*

lemma *True nitpick[satisfy] oops* — Consistency: model found

Modal collapse and Monotheism: not implied.

lemma MC : $[\forall \Phi.(\Phi \rightarrow \Box \Phi)]$ **nitpick oops** — Countermodel

lemma MT : $[\forall x y.(((\mathcal{G} x) \wedge (\mathcal{G} y)) \rightarrow (x=y))]$ **nitpick oops** — Countermodel
end

4.7 Simplified Variant with Simple Entailment in Logic K (Fig. 8 in [4])

theory *SimpleVariantSE* **imports**

HOML

MFilter

BaseDefs

begin

Axiom's of new variant based on ultrafilters.

axiomatization where

$A1'$: $[\neg(\mathcal{P}(\lambda x.(x \neq x)))]$ **and**

$A2''$: $[\forall X Y.(((\mathcal{P} X) \wedge (X \sqsubseteq Y)) \rightarrow (\mathcal{P} Y))]$ **and**

$T2$: $[\mathcal{P} \mathcal{G}]$

Necessary existence of a Godlike entity.

theorem $T6$: $[\Box(\exists^E \mathcal{G})]$ **using** $A1'$ $A2''$ $T2$ **by** *blast*

theorem $T7$: $[\exists^E \mathcal{G}]$ **using** $A1'$ $A2''$ $T2$ **by** *blast*

Possible existence of a Godlike: has counterodel.

lemma $T3$: $[\diamond(\exists^E \mathcal{G})]$ **nitpick oops** — Countermodel

lemma $T3'$: **assumes** T : $[\forall \varphi.((\Box \varphi) \rightarrow \varphi)]$

shows $[\diamond(\exists^E \mathcal{G})]$

using $A1'$ $A2''$ $T2$ T **by** *metis*

end

4.8 Simplified Variant with Simple Entailment in Logic T (Fig. 9 in [4])

theory *SimpleVariantSEinT* **imports**

HOML

MFilter

BaseDefs

begin

Axiom's of new variant based on ultrafilters.

axiomatization where

$A1'$: $[\neg(\mathcal{P}(\lambda x.(x \neq x)))]$ **and**

$A2''$: $[\forall X Y.((\mathcal{P} X) \wedge (X \sqsubseteq Y)) \rightarrow (\mathcal{P} Y)]$ and
 $T2$: $[\mathcal{P} \mathcal{G}]$

Modal Logic T.

axiomatization where T : $[\forall \varphi.((\Box \varphi) \rightarrow \varphi)]$
lemma T' : $[\forall \varphi.(\varphi \rightarrow (\Diamond \varphi))]$ **by** (*metis T*)

Necessary existence of a Godlike entity.

theorem $T6$: $[\Box(\exists^E \mathcal{G})]$ — Proof found by sledgehammer

proof –

have $T1$: $[\forall X.((\mathcal{P} X) \rightarrow (\Diamond(\exists^E X)))]$ **by** (*metis A1' A2'' T'*)
have $T3$: $[\Diamond(\exists^E \mathcal{G})]$ **by** (*metis T1 T2*)
have $T5$: $[(\Diamond(\exists^E \mathcal{G})) \rightarrow \Box(\exists^E \mathcal{G})]$ **by** (*metis A1' A2'' T2*)
thus *?thesis* **using** $T3$ **by** *simp qed*

T6 again, with an alternative, simpler proof.

theorem $T6$ *again*: $[\Box(\exists^E \mathcal{G})]$

proof –

have $L1$: $[(\exists X.((\mathcal{P} X) \wedge \neg(\exists^E X))) \rightarrow (\mathcal{P}(\lambda x.(x \neq x)))]$
by (*smt A2''*)
have $L2$: $[\neg(\exists X.((\mathcal{P} X) \wedge \neg(\exists^E X)))]$ **by** (*metis L1 A1'*)
have $T1'$: $[\forall X.((\mathcal{P} X) \rightarrow (\exists^E X))]$ **by** (*metis L2*)
have $T3'$: $[\exists^E \mathcal{G}]$ **by** (*metis T1' T2*)
have $L3$: $[\Diamond(\exists^E \mathcal{G})]$ **by** (*metis T3' T'*) — not needed
thus *?thesis* **using** $T3'$ **by** *simp qed*

end

4.9 Hauptfiltervariant (Fig. 10 in [4])

theory *SimpleVariantHF* **imports**

HOML

MFilter

BaseDefs

DisableKodkodScala

begin

Definition: Set of supersets of X , we call this $\mathcal{HF} X$.

abbreviation $HF::\gamma \Rightarrow (\gamma \Rightarrow \sigma)$ **where** $HF X \equiv \lambda Y.(X \sqsubseteq Y)$

Postulate: $\mathcal{HF} \mathcal{G}$ is a filter; i.e., Hauptfilter of \mathcal{G} .

axiomatization where $F1$: $[Filter (HF \mathcal{G})]$

Necessary existence of a Godlike entity.

theorem $T6$: $[\Box(\exists^E \mathcal{G})]$ **using** $F1$ **by** *auto* — Proof found

theorem $T6$ *again*: $[\Box(\exists^E \mathcal{G})]$

proof –

have $T3'$: $[\exists^E \mathcal{G}]$ **using** $F1$ **by** *auto*
have $T6$: $[\Box(\exists^E \mathcal{G})]$ **using** $T3'$ **by** *blast*

thus *?thesis* **by simp qed**

Possible existence of Godlike entity not implied.

lemma *T3*: $[\diamond(\exists^E \mathcal{G})]$ **nitpick oops** — Countermodel

Axiom T enforces possible existence of Godlike entity.

axiomatization

lemma *T3*: **assumes** *T*: $[\forall \varphi.((\Box \varphi) \rightarrow \varphi)]$
shows $[\diamond(\exists^E \mathcal{G})]$ **using** *F1 T* **by auto**

lemma *True* **nitpick**[*satisfy*] **oops** — Consistency: model found

Modal collapse: not implied anymore.

lemma *MC*: $[\forall \Phi.(\Phi \rightarrow \Box \Phi)]$ **nitpick oops** — Countermodel

lemma *MT*: $[\forall x y.(((\mathcal{G} x) \wedge (\mathcal{G} y)) \rightarrow (x=y))]$
nitpick oops — Countermodel

end

4.10 Formal Study of Version No.2 of Gödel's Argument as Reported by Kanckos and Lethen, 2019 [6] (Fig. 11 in [4])

theory *KanckosLethenNo2Possibilist* **imports**

HOML

MFilter

BaseDefs

begin

Axioms of Version No. 2 [6].

abbreviation *delta* (Δ) **where** $\Delta A \equiv \lambda x.(\forall \psi. ((A \psi) \rightarrow (\psi x)))$

abbreviation *N* (\mathcal{N}) **where** $\mathcal{N} \varphi \equiv \lambda x.(\Box(\varphi x))$

axiomatization where

Axiom1: $[\forall \varphi \psi.(((\mathcal{P} \varphi) \wedge (\Box(\forall x. ((\varphi x) \rightarrow (\psi x)))))) \rightarrow (\mathcal{P} \psi)]$ **and** — The \Box can be omitted here; the proofs still work.

Axiom2: $[\forall A.(\Box((\forall \varphi.((A \varphi) \rightarrow (\mathcal{P} \varphi))) \rightarrow (\mathcal{P} (\Delta A))))]$ **and** — The \Box can be omitted here; the proofs still work.

Axiom3: $[\forall \varphi.((\mathcal{P} \varphi) \rightarrow (\mathcal{P} (\mathcal{N} \varphi)))]$ **and**

Axiom4: $[\forall \varphi.((\mathcal{P} \varphi) \rightarrow (\neg(\mathcal{P}(\neg\varphi))))]$ **and**

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axB: $[\forall \varphi.(\varphi \rightarrow \Box \diamond \varphi)]$ **and** *axM*: $[\forall \varphi.((\Box \varphi) \rightarrow \varphi)]$ **and** *ax4*: $[\forall \varphi.((\Box \varphi) \rightarrow (\Box \Box \varphi))]$

Sahlqvist correspondences: they are better suited for proof automation.

lemma *axB'*: $\forall x y. \neg(xry) \vee (yrx)$ **using** *axB* **by fastforce**

lemma *axM'*: $\forall x. (xx)$ **using** *axM* **by blast**

lemma *ax4'*: $\forall x y z. (((xy) \wedge (yz)) \rightarrow (xz))$ **using** *ax4* **by auto**

Proofs for all theorems for No.2 from [6].

theorem *Theorem0*: $[\forall \varphi \psi. ((\forall Q. ((Q \varphi) \rightarrow (Q \psi))) \rightarrow ((\mathcal{P} \varphi) \rightarrow (\mathcal{P} \psi)))]$
by *auto* — not needed
theorem *Theorem1*: $[\mathcal{P} \mathcal{G}]$ **unfolding** *G-def* **using** *Axiom2 axM* **by** *blast*
theorem *Theorem2*: $[\forall x. ((\mathcal{G} x) \rightarrow (\exists y. \mathcal{G} y))]$ **by** *blast* — not needed
theorem *Theorem3a*: $[\mathcal{P} (\lambda x. (\exists y. \mathcal{G} y))]$ **by** (*metis (no-types, lifting) Axiom1 Theorem1*)
theorem *Theorem3b*: $[\Box (\mathcal{P} (\lambda x. (\Box (\exists y. \mathcal{G} y))))]$ **by** (*smt Axiom1 G-def Theorem3a Axiom3 Theorem1 axB'*)
theorem *Theorem4*: $[\forall x. \Box ((\mathcal{G} x) \rightarrow ((\mathcal{P} (\lambda x. (\Box (\exists y. \mathcal{G} y)))) \rightarrow (\Box (\exists y. \mathcal{G} y))))]$
using *G-def* **by** *fastforce* — not needed
theorem *Theorem5*: $[\forall x. \Box ((\mathcal{G} x) \rightarrow (\Box (\exists y. \mathcal{G} y)))]$ **by** (*smt (verit) G-def Theorem3a Theorem3b*) — not needed
theorem *Theorem6*: $[\Box ((\exists y. \mathcal{G} y) \rightarrow (\Box (\exists y. \mathcal{G} y)))]$ **by** (*smt G-def Theorem3a Theorem3b*)
theorem *Theorem7*: $[\Box ((\Diamond (\exists y. \mathcal{G} y)) \rightarrow (\Box (\exists y. \mathcal{G} y)))]$ **using** *Theorem6 axB'*
by *blast*
theorem *Theorem8*: $[\Box (\exists y. \mathcal{G} y)]$ **by** (*metis Axiom1 Axiom4 Theorem1 Theorem7 axB'*)
theorem *Theorem9*: $[\forall \varphi. ((\mathcal{P} \varphi) \rightarrow \Diamond (\exists x. \varphi x))]$ **using** *Axiom1 Axiom4 axM'*
by *metis*

Short proof of Theorem8; analogous to the one presented in Sec. 7 of Benzmüller 2020.

theorem $[\Box (\exists y. \mathcal{G} y)]$ — Note: this version of the proof uses only *axB'* and *axM'*.
proof —
have *L1*: $[(\exists X. ((\mathcal{P} X) \wedge \neg (\exists X))) \rightarrow (\mathcal{P} (\lambda x. (x \neq x)))]$ **using** *Axiom1 Axiom3 axB'*
by *blast* — Use *metis* here if \Box is omitted in *Axiom1* and *Axiom 2*
have *L2*: $[\neg (\mathcal{P} (\lambda x. (x \neq x)))]$ **using** *Axiom1 Axiom4* **by** *metis*
have *L3*: $[\neg (\exists X. ((\mathcal{P} X) \wedge \neg (\exists X)))]$ **using** *L1 L2* **by** *blast*
have *T2*: $[\mathcal{P} \mathcal{G}]$ **by** (*smt Axiom1 Axiom2 G-def axM'*)
have *T3*: $[\exists y. \mathcal{G} y]$ **using** *L3 T2* **by** *blast*
have *T6*: $[\Box (\exists y. \mathcal{G} y)]$ **by** (*simp add: T3*)
thus *?thesis* **by** *blast qed*

theorem *T5*: $[(\Diamond (\exists y. \mathcal{G} y)) \rightarrow \Box (\exists y. \mathcal{G} y)]$ — Obvious: If we can prove Theorem8, then we also have T5.

proof —
have *L1*: $[(\exists X. ((\mathcal{P} X) \wedge \neg (\exists X))) \rightarrow (\mathcal{P} (\lambda x. (x \neq x)))]$ **using** *Axiom1 Axiom3 axB'*
by *blast* — Use *metis* here if \Box is omitted in *Axiom1* and *Axiom 2*
have *L2*: $[\neg (\mathcal{P} (\lambda x. (x \neq x)))]$ **using** *Axiom1 Axiom4* **by** *metis*
have *L3*: $[\neg (\exists X. ((\mathcal{P} X) \wedge \neg (\exists X)))]$ **using** *L1 L2* **by** *blast*
have *T2*: $[\mathcal{P} \mathcal{G}]$ **by** (*smt Axiom1 Axiom2 G-def axM'*)
have *T3*: $[\exists y. \mathcal{G} y]$ **using** *L3 T2* **by** *blast*
have *T6*: $[\Box (\exists y. \mathcal{G} y)]$ **by** (*simp add: T3*)
thus *?thesis* **by** *blast qed*

Another short proof of Theorem8.

theorem $[\Box (\exists y. \mathcal{G} y)]$ — Note: fewer assumptions used in some cases than in [6].

proof –

have $T1$: $[\mathcal{P} \mathcal{G}]$ **unfolding** G -def **using** $Axiom2$ axM **by** $blast$
have $T3a$: $[\mathcal{P} (\lambda x. (\exists y. \mathcal{G} y))]$ **by** ($metis$ (no -types, $lifting$) $Axiom1$ $T1$)
have $T3b$: $[\Box(\mathcal{P} (\lambda x. (\Box(\exists y. \mathcal{G} y))))]$ **by** (smt $Axiom1$ G -def $T3a$ $Axiom3$ $T1$ axB')
have $T6$: $[\Box((\exists y. \mathcal{G} y) \rightarrow (\Box(\exists y. \mathcal{G} y)))]$ **by** (smt G -def $T3a$ $T3b$)
have $T7$: $[\Box((\Diamond(\exists y. \mathcal{G} y)) \rightarrow (\Box(\exists y. \mathcal{G} y)))]$ **using** $T6$ axB' **by** $blast$
thus $?thesis$ **by** (smt $Axiom1$ $Axiom4$ $T3b$ axB') **qed**

Are the axioms of the simplified versions implied?

Actualist version of the axioms.

lemma $A1'$: $[\neg(\mathcal{P}(\lambda x.(x \neq x)))]$ **using** $Theorem9$ **by** $blast$

lemma $A2'$: $[\forall X Y.(((\mathcal{P} X) \wedge ((X \sqsubseteq Y) \vee (X \Rightarrow Y))) \rightarrow (\mathcal{P} Y))]$ **nitpick oops**
— Countermodel

lemma $A3$: $[\forall \mathcal{Z}.((\mathcal{P}os \mathcal{Z}) \rightarrow (\forall X.((X \sqcap \mathcal{Z}) \rightarrow (\mathcal{P} X)))]]$ **nitpick oops** —
Countermodel

Possibilist version of the axioms.

abbreviation a ($-\sqsubseteq^p$ -) **where** $X \sqsubseteq^p Y \equiv \forall z.((X z) \rightarrow (Y z))$

abbreviation b ($-\Rightarrow^p$ -) **where** $X \Rightarrow^p Y \equiv \Box(X \sqsubseteq^p Y)$

abbreviation d ($-\sqcap^p$ -) **where** $X \sqcap^p \mathcal{Z} \equiv \Box(\forall u.((X u) \leftrightarrow (\forall Y.((\mathcal{Z} Y) \rightarrow (Y u)))))$

lemma $A1'P$: $[\neg(\mathcal{P}(\lambda x.(x \neq x)))]$ **using** $Theorem9$ **by** $blast$

lemma $A2'P$: $[\forall X Y.(((\mathcal{P} X) \wedge ((X \sqsubseteq^p Y) \vee (X \Rightarrow^p Y))) \rightarrow (\mathcal{P} Y))]$ **oops** — no
answer, yet by sledgehammer and nitpick

lemma $A2'aP$: $[\forall X Y.(((\mathcal{P} X) \wedge (X \Rightarrow^p Y)) \rightarrow (\mathcal{P} Y))]$ **using** $Axiom1$ axM' **by**
 $metis$

lemma $A2'bP$: $[\forall X Y.(((\mathcal{P} X) \wedge (X \sqsubseteq^p Y)) \rightarrow (\mathcal{P} Y))]$ **oops** — no answer, yet
by sledgehammer and nitpick

lemma $A3P$: $[\forall \mathcal{Z}.((\mathcal{P}os \mathcal{Z}) \rightarrow (\forall X.((X \sqcap^p \mathcal{Z}) \rightarrow (\mathcal{P} X)))]]$

by (smt ($verit$, del -insts) $Axiom1$ $Axiom2$ axM') — proof found

Are Axiom2 and A3 equivalent? Only when assuming Axiom1 and axiom
M.

lemma $[\forall A.(\Box((\forall \varphi.((A \varphi) \rightarrow (\mathcal{P} \varphi)) \rightarrow (\mathcal{P} (\Delta A)))))] \equiv [\forall \mathcal{Z}.((\mathcal{P}os \mathcal{Z}) \rightarrow$
 $(\forall X.((X \sqcap^p \mathcal{Z}) \rightarrow (\mathcal{P} X)))]]$

by (smt ($verit$, $ccfv$ -threshold) $Axiom1$ axM') — proof found
end

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