

A Formally Verified Checker of the Safe Distance Traffic Rules for Autonomous Vehicles

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March 17, 2025

Abstract

The Vienna Convention on Road Traffic defines the safe distance traffic rules informally. This could make autonomous vehicle liable for safe-distance-related accidents because there is no clear definition of how large a safe distance is. We provide a formally proven prescriptive definition of a safe distance, and checkers which can decide whether an autonomous vehicle is obeying the safe distance rule. Not only does our work apply to the domain of law, but it also serves as a specification for autonomous vehicle manufacturers and for online verification of path planners. This formalization accompanies our paper "A Formally Verified Checker of the Safe Distance Traffic Rules for Autonomous Vehicles". [1]

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1 Safe Distance

```
theory Safe-Distance
imports
  HOL-Analysis.Multivariate-Analysis
  HOL-Decision-Proc.Approximation
  Sturm-Sequences.Sturm
begin
```

This theory is about formalising the safe distance rule. The safe distance rule is obtained from Vienna Convention which basically states the following thing.

“The car at all times must maintain a safe distance towards the vehicle in front of it, such that whenever the vehicle in front and the ego vehicle apply maximum deceleration, there will not be a collision.”

To formalise this safe distance rule we have to define first what is a safe distance. To define this safe distance, we have to model the physics of the movement of the vehicle. The following model is sufficient.

$$s = s_0 + v_0 * t + 1 / 2 * a_0 * t^2$$

Assumptions in this model are :

- Both vehicles are assumed to be point mass. The exact location of the ego vehicle is the front-most occupancy of the ego vehicle. Similarly for the other vehicle, its exact location is the rearmost occupancy of the other vehicle.
- Both cars can never drive backward.

```
lemmas [simp del] = div-mult-self1 div-mult-self2 div-mult-self3 div-mult-self4
```

1.1 Quadratic Equations

```
lemma discriminant:  $a * x^2 + b * x + c = (0::real) \implies 0 \leq b^2 - 4 * a * c$ 
  ⟨proof⟩
```

```
lemma quadratic-eq-factoring:
```

```

assumes D : D = b2 - 4 * a * c
assumes nn: 0 ≤ D
assumes x1: x1 = (-b + sqrt D) / (2 * a)
assumes x2: x2 = (-b - sqrt D) / (2 * a)
assumes a : a ≠ 0
shows a * x2 + b * x + c = a * (x - x1) * (x - x2)
⟨proof⟩

lemma quadratic-eq-zeroes-iff:
assumes D : D = b2 - 4 * a * c
assumes x1: x1 = (-b + sqrt D) / (2 * a)
assumes x2: x2 = (-b - sqrt D) / (2 * a)
assumes a : a ≠ 0
shows a * x2 + b * x + c = 0 ↔ (D ≥ 0 ∧ (x = x1 ∨ x = x2)) (is ?z ↔ -)
⟨proof⟩

```

1.2 Convexity Condition

```

lemma p-convex:
fixes a b c x y z :: real
assumes p-def: p = (λx. a * x2 + b * x + c)
assumes less : x < y y < z
and ge : p x > p y p y ≤ p z
shows a > 0
⟨proof⟩

```

```

definition root-in :: real ⇒ real ⇒ (real ⇒ real) ⇒ bool where
root-in m M f = (exists x ∈ {m .. M}. f x = 0)

```

```

definition quadroot-in :: real ⇒ real ⇒ real ⇒ real ⇒ real ⇒ bool where
quadroot-in m M a b c = root-in m M (λx. a * x2 + b * x + c)

```

```

lemma card-iff-exists: 0 < card X ↔ finite X ∧ (exists x. x ∈ X)
⟨proof⟩

```

```

lemma quadroot-in-sturm[code]:
quadroot-in m M a b c ↔ (a = 0 ∧ b = 0 ∧ c = 0 ∧ m ≤ M) ∨
(m ≤ M ∧ poly [:c, b, a:] m = 0) ∨
count-roots-between [:c, b, a:] m M > 0
⟨proof⟩

```

```

lemma check-quadroot-linear:
fixes a b c :: real
assumes a = 0
shows ¬ quadroot-in m M a b c ↔
((b = 0 ∧ c = 0 ∧ M < m) ∨ (b = 0 ∧ c ≠ 0) ∨
(b ≠ 0 ∧ (let x = -c / b in m > x ∨ x > M)))
⟨proof⟩

```

```

lemma check-quadroot-nonlinear:
  assumes a ≠ 0
  shows quadroot-in m M a b c =
    (let D = b ^ 2 - 4 * a * c in D ≥ 0 ∧
     ((let x = (-b + sqrt D)/(2*a) in m ≤ x ∧ x ≤ M) ∨
      (let x = (-b - sqrt D)/(2*a) in m ≤ x ∧ x ≤ M)))
  ⟨proof⟩

lemma ncheck-quadroot:
  shows ¬quadroot-in m M a b c ↔
    (a = 0 → ¬quadroot-in m M a b c) ∧
    (a = 0 ∨ ¬quadroot-in m M a b c)
  ⟨proof⟩

```

1.3 Movement

```

locale movement =
  fixes a v s0 :: real
begin

```

Function to compute the distance using the equation

$$s(t) = s_0 + v_0 * t + 1 / 2 * a_0 * t^2$$

Input parameters:

- s_0 : initial distance
- v_0 : initial velocity (positive means forward direction and the converse is true)
- a : acceleration (positive for increasing and negative for decreasing)
- t : time

For the time $t < 0$, we assume the output of the function is s_0 . Otherwise, the output is calculated according to the equation above.

1.3.1 Continuous Dynamics

```

definition p :: real ⇒ real where
  p t = s0 + v * t + 1/2 * a * t^2

```

```

lemma p-all-zeroes:
  assumes D: D = v^2 - 2 * a * s0
  shows p t = 0 ↔ ((a ≠ 0 ∧ 0 ≤ D ∧ ((t = (-v + sqrt D) / a) ∨ t = (-v - sqrt D) / a)) ∨
    (a = 0 ∧ v = 0 ∧ s0 = 0) ∨ (a = 0 ∧ v ≠ 0 ∧ t = (-s0 / v)))
  ⟨proof⟩

```

lemma $p\text{-zero}[simp]$: $p\ 0 = s0$
 $\langle proof \rangle$

lemma $p\text{-continuous}[continuous\text{-intros}]$: continuous-on $T\ p$
 $\langle proof \rangle$

lemma $isCont\ p[p[continuous\text{-intros}]]$: $isCont\ p\ x$
 $\langle proof \rangle$

definition $p' :: real \Rightarrow real$ **where**

$$p' t = v + a * t$$

lemma $p'\text{-zero}$: $p' 0 = v$
 $\langle proof \rangle$

lemma $p\text{-has-vector-derivative}[derivative\text{-intros}]$: (p has-vector-derivative $p' t$) (at t within s)
 $\langle proof \rangle$

lemma $p\text{-has-real-derivative}[derivative\text{-intros}]$: (p has-real-derivative $p' t$) (at t within s)
 $\langle proof \rangle$

definition $p'' :: real \Rightarrow real$ **where**

$$p'' t = a$$

lemma $p'\text{-has-vector-derivative}[derivative\text{-intros}]$: (p' has-vector-derivative $p'' t$) (at t within s)
 $\langle proof \rangle$

lemma $p'\text{-has-real-derivative}[derivative\text{-intros}]$: (p' has-real-derivative $p'' t$) (at t within s)
 $\langle proof \rangle$

definition $t\text{-stop} :: real$ **where**

$$t\text{-stop} = -v / a$$

lemma $p'\text{-stop-zero}$: $p' t\text{-stop} = (\text{if } a = 0 \text{ then } v \text{ else } 0)$ $\langle proof \rangle$

lemma $p'\text{-pos-iff}$: $p' x > 0 \longleftrightarrow (\text{if } a > 0 \text{ then } x > -v / a \text{ else if } a < 0 \text{ then } x < -v / a \text{ else } v > 0)$
 $\langle proof \rangle$

lemma $le\text{-}t\text{-stop-iff}$: $a \neq 0 \implies x \leq t\text{-stop} \longleftrightarrow (\text{if } a < 0 \text{ then } p' x \geq 0 \text{ else } p' x \leq 0)$
 $\langle proof \rangle$

lemma $p'\text{-continuous}[continuous\text{-intros}]$: continuous-on $T\ p'$
 $\langle proof \rangle$

```
lemma isCont-p[continuous-intros]: isCont p' x
  ⟨proof⟩
```

```
definition p-max :: real where
  p-max = p t-stop
```

```
lemmas p-t-stop = p-max-def[symmetric]
```

```
lemma p-max-eq: p-max = s0 - v^2 / a / 2
  ⟨proof⟩
```

1.3.2 Hybrid Dynamics

```
definition s :: real ⇒ real where
  s t = ( if t ≤ 0 then s0
            else if t ≤ t-stop then p t
            else p-max)
```

```
definition q :: real ⇒ real where
  q t = s0 + v * t
```

```
definition q' :: real ⇒ real where
  q' t = v
```

```
lemma init-q: q 0 = s0 ⟨proof⟩
```

```
lemma q-continuous[continuous-intros]: continuous-on T q
  ⟨proof⟩
```

```
lemma isCont-q[continuous-intros]: isCont q x
  ⟨proof⟩
```

```
lemma q-has-vector-derivative[derivative-intros]: (q has-vector-derivative q' t) (at t within u)
  ⟨proof⟩
```

```
lemma q-has-real-derivative[derivative-intros]: (q has-real-derivative q' t) (at t within u)
  ⟨proof⟩
```

```
lemma s-cond-def:
  t ≤ 0 ⇒⇒ s t = s0
  0 ≤ t ≤ t-stop ⇒⇒ s t = p t
  ⟨proof⟩
```

```
end
```

```
locale braking-movement = movement +
```

```

assumes decel:  $a < 0$ 
assumes nonneg-vel:  $v \geq 0$ 
begin

lemma t-stop-nonneg:  $0 \leq t\text{-stop}$ 
  ⟨proof⟩

lemma t-stop-pos:
  assumes  $v \neq 0$ 
  shows  $0 < t\text{-stop}$ 
  ⟨proof⟩

lemma t-stop-zero:
  assumes  $t\text{-stop} = 0$ 
  shows  $v = 0$ 
  ⟨proof⟩

lemma t-stop-zero-not-moving:  $t\text{-stop} = 0 \implies q t = s0$ 
  ⟨proof⟩

abbreviation s-stop ≡ s t-stop

lemma s-t-stop:  $s\text{-stop} = p\text{-max}$ 
  ⟨proof⟩

lemma s0-le-s-stop:  $s0 \leq s\text{-stop}$ 
  ⟨proof⟩

lemma p-mono:  $x \leq y \implies y \leq t\text{-stop} \implies p x \leq p y$ 
  ⟨proof⟩

lemma p-antimono:  $x \leq y \implies t\text{-stop} \leq x \implies p y \leq p x$ 
  ⟨proof⟩

lemma p-strict-mono:  $x < y \implies y \leq t\text{-stop} \implies p x < p y$ 
  ⟨proof⟩

lemma p-strict-antimono:  $x < y \implies t\text{-stop} \leq x \implies p y < p x$ 
  ⟨proof⟩

lemma p-max:  $p x \leq p\text{-max}$ 
  ⟨proof⟩

lemma continuous-on-s[continuous-intros]: continuous-on T s
  ⟨proof⟩

lemma isCont-s[continuous-intros]: isCont s x
  ⟨proof⟩

```

```

definition s' :: real  $\Rightarrow$  real where
  s' t = (if  $t \leq t\text{-stop}$  then p' t else 0)

lemma s-has-real-derivative:
  assumes  $t \geq 0 v / a \leq 0 a \neq 0$ 
  shows (s has-real-derivative s' t) (at t within {0..})
  ⟨proof⟩

lemma s-has-vector-derivative[derivative-intros]:
  assumes  $t \geq 0 v / a \leq 0 a \neq 0$ 
  shows (s has-vector-derivative s' t) (at t within {0..})
  ⟨proof⟩

lemma s-has-field-derivative[derivative-intros]:
  assumes  $t \geq 0 v / a \leq 0 a \neq 0$ 
  shows (s has-field-derivative s' t) (at t within {0..})
  ⟨proof⟩

lemma s-has-real-derivative-at:
  assumes  $0 < x 0 \leq v a < 0$ 
  shows (s has-real-derivative s' x) (at x)
  ⟨proof⟩

lemma s-delayed-has-field-derivative[derivative-intros]:
  assumes  $\delta < t 0 \leq v a < 0$ 
  shows (( $\lambda x. s(x - \delta)$ ) has-field-derivative s' (t -  $\delta$ )) (at t within { $\delta < ..$ })
  ⟨proof⟩

lemma s-delayed-has-vector-derivative[derivative-intros]:
  assumes  $\delta < t 0 \leq v a < 0$ 
  shows (( $\lambda x. s(x - \delta)$ ) has-vector-derivative s' (t -  $\delta$ )) (at t within { $\delta < ..$ })
  ⟨proof⟩

lemma s'-nonneg:  $0 \leq v \implies a \leq 0 \implies 0 \leq s' x$ 
  ⟨proof⟩

lemma s'-pos:  $0 \leq x \implies x < t\text{-stop} \implies 0 \leq v \implies a \leq 0 \implies 0 < s' x$ 
  ⟨proof⟩

```

1.3.3 Monotonicity of Movement

```

lemma s-mono:
  assumes  $t \geq u u \geq 0$ 
  shows s t  $\geq$  s u
  ⟨proof⟩

lemma s-strict-mono:
  assumes  $u < t t \leq t\text{-stop} u \geq 0$ 
  shows s u  $<$  s t

```

$\langle proof \rangle$

```
lemma s-antimono:  
  assumes x ≤ y  
  assumes t-stop ≤ x  
  shows s y ≤ s x  
 $\langle proof \rangle$ 
```

```
lemma init-s : t ≤ 0 ⇒ s t = s0  
 $\langle proof \rangle$ 
```

```
lemma q-min: 0 ≤ t ⇒ s0 ≤ q t  
 $\langle proof \rangle$ 
```

```
lemma q-mono: x ≤ y ⇒ q x ≤ q y  
 $\langle proof \rangle$ 
```

1.3.4 Maximum at Stopping Time

```
lemma s-max: s x ≤ s-stop  
 $\langle proof \rangle$ 
```

```
lemma s-eq-s-stop: NO-MATCH t-stop x ⇒ x ≥ t-stop ⇒ s x = s-stop  
 $\langle proof \rangle$ 
```

end

1.4 Safe Distance

```
locale safe-distance =  
  fixes ae ve se :: real  
  fixes ao vo so :: real  
  assumes nonneg-vel-ego : 0 ≤ ve  
  assumes nonneg-vel-other : 0 ≤ vo  
  assumes decelerate-ego : ae < 0  
  assumes decelerate-other : ao < 0  
  assumes in-front : se < so  
begin
```

```
lemmas hyps =  
  nonneg-vel-ego  
  nonneg-vel-other  
  decelerate-ego  
  decelerate-other  
  in-front
```

```
sublocale ego: braking-movement ae ve se  $\langle proof \rangle$   
sublocale other: braking-movement ao vo so  $\langle proof \rangle$   
sublocale ego-other: movement ao − ae vo − ve so − se  $\langle proof \rangle$ 
```

1.4.1 Collision

definition *collision* :: *real set* \Rightarrow *bool* **where**

$$\text{collision time-set} \equiv (\exists t \in \text{time-set}. \text{ego.s } t = \text{other.s } t)$$

abbreviation *no-collision* :: *real set* \Rightarrow *bool* **where**

$$\text{no-collision time-set} \equiv \neg \text{collision time-set}$$

lemma *no-collision-initially* : *no-collision* {.. 0}

{proof}

lemma *no-collisionI*:

$$(\bigwedge t. t \in S \implies \text{ego.s } t \neq \text{other.s } t) \implies \text{no-collision } S$$

{proof}

theorem *cond-1*: *ego.s-stop* $<$ *s_o* \implies *no-collision* {0..}

{proof}

lemma *ego-other-strict-intv*:

assumes *ego.s t* $>$ *other.s t*
shows *collision* {0 .. < *t*}

{proof}

lemma *collision-subset*: *collision s* \implies *s* \subseteq *t* \implies *collision t*

{proof}

lemma *ego-other-intv*:

assumes *ego.s t* \geq *other.s t*
shows *collision* {0 .. *t*}
{proof}

theorem *cond-2*:

assumes *ego.s-stop* \geq *other.s-stop*
shows *collision* {0 ..}
{proof}

abbreviation *D2* :: *real* **where**

$$D2 \equiv (v_o - v_e)^2 - 2 * (a_o - a_e) * (s_o - s_e)$$

abbreviation *tD'* :: *real* **where**

$$tD' \equiv \sqrt{2 * (\text{ego.s-stop} - \text{other.s-stop}) / a_o}$$

lemma *pos-via-half-dist*:

dist a b $<$ *b* / 2 \implies *b* $>$ 0 \implies *a* $>$ 0
{proof}

lemma *collision-within-p*:

assumes *s_o* \leq *ego.s-stop* *ego.s-stop* $<$ *other.s-stop*
shows *collision* {0..} \longleftrightarrow ($\exists t \geq 0.$ *ego.p t* = *other.p t* \wedge *t* $<$ *ego.t-stop* \wedge *t* $<$ *other.t-stop*)

$\langle proof \rangle$

lemma *collision-within-eq*:

assumes $s_o \leq ego.s\text{-stop}$ $ego.s\text{-stop} < other.s\text{-stop}$
shows $collision \{0..\} \longleftrightarrow collision \{0 .. < \min ego.t\text{-stop} other.t\text{-stop}\}$
 $\langle proof \rangle$

lemma *collision-excluded*: $(\bigwedge t. t \in T \implies ego.s t \neq other.s t) \implies collision S \longleftrightarrow collision (S - T)$

$\langle proof \rangle$

lemma *collision-within-less*:

assumes $s_o \leq ego.s\text{-stop}$ $ego.s\text{-stop} < other.s\text{-stop}$
shows $collision \{0..\} \longleftrightarrow collision \{0 .. < \min ego.t\text{-stop} other.t\text{-stop}\}$
 $\langle proof \rangle$

theorem *cond-3*:

assumes $s_o \leq ego.s\text{-stop}$ $ego.s\text{-stop} < other.s\text{-stop}$
shows $collision \{0..\} \longleftrightarrow (a_o > a_e \wedge v_o < v_e \wedge 0 \leq D2 \wedge \sqrt{D2} > v_e - a_e / a_o * v_o)$
 $\langle proof \rangle$

1.4.2 Formalising Safe Distance

First definition for Safe Distance based on *cond-1*.

definition *absolute-safe-distance* :: real **where**
 $absolute-safe-distance = -v_e^2 / (2 * a_e)$

lemma *absolute-safe-distance*:

assumes $s_o - s_e > absolute-safe-distance$
shows *no-collision* $\{0..\}$
 $\langle proof \rangle$

First Fallback for Safe Distance.

definition *fst-safe-distance* :: real **where**
 $fst-safe-distance = v_o^2 / (2 * a_o) - v_e^2 / (2 * a_e)$

definition *distance-leq-d2* :: real **where**
 $distance-leq-d2 = (a_e + a_o) / (2 * a_o^2) * v_o^2 - v_o * v_e / a_o$

lemma *snd-leq-fst-exp*: $distance-leq-d2 \leq fst-safe-distance$
 $\langle proof \rangle$

lemma *sqrt-D2-leq-stop-time-diff*:

assumes $a_e < a_o$
assumes $0 \leq v_e - a_e / a_o * v_o$
assumes $s_o - s_e \geq distance-leq-d2$
shows $\sqrt{D2} \leq v_e - a_e / a_o * v_o$
 $\langle proof \rangle$

```

lemma cond2-imp-pos-vo:
  assumes  $s_o \leq ego.s\text{-stop}$   $ego.s\text{-stop} < other.s\text{-stop}$ 
  shows  $v_o \neq 0$ 
  (proof)

lemma cond2-imp-gt-fst-sd:
  assumes  $s_o \leq ego.s\text{-stop}$   $ego.s\text{-stop} < other.s\text{-stop}$ 
  shows  $fst\text{-safe-distance} < s_o - s_e$ 
  (proof)

Second Fallback for Safe Distance.

definition snd-safe-distance :: real where
   $snd\text{-safe-distance} = (v_o - v_e)^2 / (2 * (a_o - a_e))$ 

lemma fst-leq-snd-safe-distance:
  assumes  $a_e < a_o$ 
  shows  $fst\text{-safe-distance} \leq snd\text{-safe-distance}$ 
  (proof)

lemma snd-safe-distance-iff-nonneg-D2:
  assumes  $a_e < a_o$ 
  shows  $s_o - s_e \leq snd\text{-safe-distance} \longleftrightarrow 0 \leq D2$ 
  (proof)

lemma t-stop-diff-neg-means-leq-D2:
  assumes  $s_o \leq ego.s\text{-stop}$   $ego.s\text{-stop} < other.s\text{-stop}$   $a_e < a_o$   $0 \leq D2$ 
  shows  $v_e - a_e / a_o * v_o < 0 \longleftrightarrow \sqrt{D2} > v_e - a_e / a_o * v_o$ 
  (proof)

theorem cond-3':
  assumes  $s_o \leq ego.s\text{-stop}$   $ego.s\text{-stop} < other.s\text{-stop}$ 
  shows  $collision \{0..\} \longleftrightarrow (a_o > a_e \wedge v_o < v_e \wedge s_o - s_e \leq snd\text{-safe-distance} \wedge$ 
 $v_e - a_e / a_o * v_o < 0)$ 
  (proof)

definition d :: real  $\Rightarrow$  real where
   $d t = ($ 
    if  $t \leq 0$  then  $s_o - s_e$ 
    else if  $t \leq ego.t\text{-stop} \wedge t \leq other.t\text{-stop}$  then  $ego\text{-other}.p t$ 
    else if  $ego.t\text{-stop} \leq t \wedge t \leq other.t\text{-stop}$  then  $other.p t - ego.s\text{-stop}$ 
    else if  $other.t\text{-stop} \leq t \wedge t \leq ego.t\text{-stop}$  then  $other.s\text{-stop} - ego.p t$ 
    else  $other.s\text{-stop} - ego.s\text{-stop}$ 
   $)$ 

lemma d-diff:  $d t = other.s t - ego.s t$ 
  (proof)

lemma collision-d:  $collision S \longleftrightarrow (\exists t \in S. d t = 0)$ 

```

$\langle proof \rangle$

lemma *collision-restrict*: *collision* {0..} \longleftrightarrow *collision* {0..max *ego.t-stop* *other.t-stop*}

$\langle proof \rangle$

lemma *collision-union*: *collision* (*A* \cup *B*) \longleftrightarrow *collision A* \vee *collision B*

$\langle proof \rangle$

lemma *symbolic-checker*:

collision {0..} \longleftrightarrow

(quadroot-in 0 (min *ego.t-stop* *other.t-stop*) (1/2 * (*a_o* - *a_e*)) (*v_o* - *v_e*) (*s_o* - *s_e*)) \vee

(quadroot-in *ego.t-stop* *other.t-stop* (1/2 * *a_o*) *v_o* (*s_o* - *ego.s-stop*)) \vee

(quadroot-in *other.t-stop* *ego.t-stop* (1/2 * *a_e*) *v_e* (*s_e* - *other.s-stop*))

(is - \longleftrightarrow ?q1 \vee ?q2 \vee ?q3)

$\langle proof \rangle$

end

1.5 Checker Design

definition *rel-dist-to-stop* :: real \Rightarrow real \Rightarrow real **where**

rel-dist-to-stop *v a* \equiv - *v*² / (2 * *a*)

context includes *floatarith-syntax* **begin**

definition *rel-dist-to-stop-expr* :: nat \Rightarrow nat \Rightarrow floatarith **where**

rel-dist-to-stop-expr *v a* = *Mult* (*Minus* (*Power* (*Var v*) 2)) (*Inverse* (*Mult* (Num 2) (*Var a*)))

definition *rel-dist-to-stop'* :: nat \Rightarrow float interval option \Rightarrow float interval option
 \Rightarrow float interval option **where**

rel-dist-to-stop' *p v a* = *approx p* (*rel-dist-to-stop-expr* 0 1) [*v, a*]

lemma *rel-dist-to-stop'*: *interpret-floatarith* (*rel-dist-to-stop-expr* 0 1) [*v, a*] = *rel-dist-to-stop*

v a

$\langle proof \rangle$

definition *first-safe-dist* :: real \Rightarrow real \Rightarrow real **where**

first-safe-dist *v_e a_e* \equiv *rel-dist-to-stop* *v_e a_e*

definition *second-safe-dist* :: real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real **where**

second-safe-dist *v_e a_e v_o a_o* \equiv *rel-dist-to-stop* *v_e a_e* - *rel-dist-to-stop* *v_o a_o*

definition *second-safe-dist-expr* :: nat \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow floatarith **where**

second-safe-dist-expr *ve ae vo ao* = *Add* (*rel-dist-to-stop-expr* *ve ae*) (*Minus* (*rel-dist-to-stop-expr* *vo ao*))

definition *second-safe-dist'* :: nat \Rightarrow float interval option \Rightarrow float interval option

$\Rightarrow \text{float interval option} \Rightarrow \text{float interval option} \Rightarrow \text{float interval option}$ **where**
 $\text{second-safe-dist}' p v_e a_e v_o a_o = \text{approx } p (\text{second-safe-dist-expr } 0 1 2 3) [v_e, a_e, v_o, a_o]$

lemma *second-safe-dist'*:

*interpret-floatarith (second-safe-dist-expr 0 1 2 3) [v, a, v', a'] = second-safe-dist
 $v a v' a'$*
 $\langle proof \rangle$

definition *t-stop :: real \Rightarrow real \Rightarrow real* **where**
 $t\text{-stop } v a \equiv -v / a$

definition *t-stop-expr :: nat \Rightarrow nat \Rightarrow floatarith* **where**
 $t\text{-stop-expr } v a = \text{Minus} (\text{Mult} (\text{Var } v) (\text{Inverse} (\text{Var } a)))$
end

definition *s-stop :: real \Rightarrow real \Rightarrow real \Rightarrow real* **where**
 $s\text{-stop } s v a \equiv s + \text{rel-dist-to-stop } v a$

definition *discriminant :: real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real* **where**
 $\text{discriminant } s_e v_e a_e s_o v_o a_o \equiv (v_o - v_e)^2 - 2 * (a_o - a_e) * (s_o - s_e)$

definition *suff-cond-safe-dist2 :: real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow bool*
where
 $\text{suff-cond-safe-dist2 } s_e v_e a_e s_o v_o a_o \equiv$
 $\text{let } D2 = \text{discriminant } s_e v_e a_e s_o v_o a_o$
 $\text{in } \neg(a_e < a_o \wedge v_o < v_e \wedge 0 \leq D2 \wedge v_e - a_e / a_o * v_o < \sqrt{D2})$
 $)$

lemma *less-sqrt-iff: $y \geq 0 \implies x < \sqrt{y} \longleftrightarrow (x \geq 0 \rightarrow x^2 < y)$*
 $\langle proof \rangle$

lemma *suff-cond-safe-dist2-code[code]:*
 $\text{suff-cond-safe-dist2 } s_e v_e a_e s_o v_o a_o =$
 $(\text{let } D2 = \text{discriminant } s_e v_e a_e s_o v_o a_o \text{ in}$
 $(a_e < a_o \rightarrow v_o < v_e \rightarrow 0 \leq D2 \rightarrow (v_e - a_e / a_o * v_o \geq 0 \wedge (v_e - a_e / a_o * v_o)^2 \geq D2)))$
 $\langle proof \rangle$

There are two expressions for safe distance. The first safe distance *first-safe-dist* is always valid. Whenever the distance is bigger than *first-safe-dist*, it is guaranteed to be collision free. The second one is *second-safe-dist*. If the sufficient condition *suff-cond-safe-dist2* is satisfied and the distance is bigger than *second-safe-dist*, it is guaranteed to be collision free.

definition *check-precond :: real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow bool*
where

check-precond $s_e v_e a_e s_o v_o a_o \longleftrightarrow s_o > s_e \wedge 0 \leq v_e \wedge 0 \leq v_o \wedge a_e < 0 \wedge a_o < 0$

lemma *check-precond-safe-distance*:

check-precond $s_e \ v_e \ a_e \ s_o \ v_o \ a_o = \text{safe-distance} \ a_e \ v_e \ s_e \ a_o \ v_o \ s_o$
 $\langle \text{proof} \rangle$

1.5.1 Prescriptive Checker

definition *checker* :: $\text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{bool}$ **where**

checker $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \equiv$
 $\text{let } \text{distance} = s_o - s_e;$
 $\text{precond} = \text{check-precond} \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o;$
 $\text{safe-dist1} = \text{first-safe-dist} \ v_e \ a_e;$
 $\text{safe-dist2} = \text{second-safe-dist} \ v_e \ a_e \ v_o \ a_o;$
 $\text{cond2} = \text{suff-cond-safe-dist2} \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o$
 $\text{in precond} \wedge (\text{safe-dist1} < \text{distance} \vee (\text{safe-dist2} < \text{distance} \wedge \text{distance} \leq \text{safe-dist1} \wedge \text{cond2}))$

lemma *aux-logic*:

assumes $a \implies b$
assumes $b \implies a \longleftrightarrow c$
shows $a \longleftrightarrow b \wedge c$
 $\langle \text{proof} \rangle$

theorem *soundness-correctness*:

checker $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \longleftrightarrow \text{check-precond} \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o \wedge \text{safe-distance.no-collision}$
 $a_e \ v_e \ s_e \ a_o \ v_o \ s_o \ \{0..\}$
 $\langle \text{proof} \rangle$

definition *checker2* :: $\text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{bool}$ **where**

checker2 $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \equiv$
 $\text{let } \text{distance} = s_o - s_e;$
 $\text{precond} = \text{check-precond} \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o;$
 $\text{safe-dist1} = \text{first-safe-dist} \ v_e \ a_e;$
 $\text{safe-dist2} = \text{second-safe-dist} \ v_e \ a_e \ v_o \ a_o;$
 $\text{safe-dist3} = -\text{rel-dist-to-stop} \ (v_o - v_e) \ (a_o - a_e)$
 in
 $\text{if } \neg \text{precond} \text{ then False}$
 $\text{else if } \text{distance} > \text{safe-dist1} \text{ then True}$
 $\text{else if } a_o > a_e \wedge v_o < v_e \wedge v_e - a_e / a_o * v_o < 0 \text{ then distance} > \text{safe-dist3}$
 $\text{else distance} > \text{safe-dist2}$

theorem *checker-eq-checker2*: *checker* $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \longleftrightarrow \text{checker2} \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o$
 $\langle \text{proof} \rangle$

definition *checker3* :: $\text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{bool}$ **where**

checker3 $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \equiv$
 $\text{let } \text{distance} = s_o - s_e;$
 $\text{precond} = \text{check-precond} \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o;$
 $\text{s-stop-e} = s_e + \text{rel-dist-to-stop} \ v_e \ a_e;$

```

 $s\text{-stop-}o = s_o + \text{rel-dist-to-stop } v_o \ a_o$ 
in precond  $\wedge$  ( $s\text{-stop-}e < s_o$ 
 $\vee (s_o \leq s\text{-stop-}e \wedge s\text{-stop-}e < s\text{-stop-}o \wedge$ 
 $(\neg(a_o > a_e \wedge v_o < v_e \wedge v_e - a_e / a_o * v_o < 0 \wedge \text{distance} * (a_o -$ 
 $a_e) \leq (v_o - v_e)^2 / 2)))$ )

```

theorem checker2-eq-checker3:

```

checker2  $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \longleftrightarrow \text{checker3 } s_e \ v_e \ a_e \ s_o \ v_o \ a_o$ 
⟨proof⟩

```

1.5.2 Approximate Checker

```

lemma checker2-def': checker2  $a \ b \ c \ d \ e \ f = ($ 
let distance  $= d - a;$ 
precond  $= \text{check-precond } a \ b \ c \ d \ e \ f;$ 
safe-dist1  $= \text{first-safe-dist } b \ c;$ 
safe-dist2  $= \text{second-safe-dist } b \ c \ e \ f;$ 
C  $= c < f \wedge e < b \wedge b * f > c * e;$ 
P1  $= (e - b)^2 < 2 * \text{distance} * (f - c);$ 
P2  $= -b^2 / c + e^2 / f < 2 * \text{distance}$ 
in precond  $\wedge$  (safe-dist1  $< \text{distance} \vee$ 
safe-dist1  $\geq \text{distance} \wedge (C \wedge P1 \vee \neg C \wedge P2))$ )
⟨proof⟩

```

```

lemma power2-less-sqrt-iff:  $(x::\text{real})^2 < y \longleftrightarrow (y \geq 0 \wedge \text{abs } x < \sqrt{y})$ 
⟨proof⟩

```

```

schematic-goal checker-form: interpret-form ?x ?y  $\implies \text{checker } s_e \ v_e \ a_e \ s_o \ v_o \ a_o$ 
⟨proof⟩

```

```

context includes floatarith-syntax
begin
definition checker' p  $s_e \ v_e \ a_e \ s_o \ v_o \ a_o = \text{approx-form } p$ 
 $(\text{Conj} (\text{Conj} (\text{Less} (\text{Var} (\text{Suc} (\text{Suc} 0))) (\text{Var} (\text{Suc} (\text{Suc} 0)))))$ 
 $(\text{Conj} (\text{LessEqual} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))))$ 
 $(\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0)))))))$ 
 $(\text{Conj} (\text{LessEqual} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))))$ 
 $(\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0)))))))$ 
 $(\text{Conj} (\text{Less} (\text{Var} 0) (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))))$ 
 $(\text{Less} (\text{Var} (\text{Suc} 0)) (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))))))$ 
 $(\text{Disj} (\text{Less} (\text{Add} (\text{Var} (\text{Suc} (\text{Suc} 0)))))$ 
 $(\text{Mult} (\text{Minus} (\text{Power} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))))$ 
 $2)) (\text{Inverse} (\text{Mult} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0)))))) (\text{Var} 0))))$ 
 $(\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))$ 
 $(\text{Conj} (\text{LessEqual} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))$ 
 $(\text{Add} (\text{Var} (\text{Suc} (\text{Suc} 0))))$ 
 $(\text{Mult} (\text{Minus} (\text{Power} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} 0))))))))$ 
 $2)) (\text{Inverse} (\text{Mult} (\text{Var} (\text{Suc} (\text{Suc} (\text{Suc} 0)))))) (\text{Var} 0))))$ 

```

```

(Conj (Less (Add (Var (Suc (Suc 0))))
              (Mult (Minus (Power (Var (Suc (Suc (Suc (Suc (Suc 0)))))) 2)) (Inverse (Mult (Var (Suc (Suc (Suc (Suc 0)))))) (Var 0)))))
              (Add (Var (Suc (Suc (Suc 0)))))
              (Mult (Minus (Power (Var (Suc (Suc (Suc (Suc (Suc 0)))))) 2)))
                     (Inverse (Mult (Var (Suc (Suc (Suc (Suc 0)))))) (Var (Suc 0)))))))
              (Disj (LessEqual (Var (Suc 0)) (Var 0))
                     (Disj (LessEqual (Var (Suc (Suc (Suc (Suc 0)))))) (Var (Suc (Suc (Suc (Suc (Suc 0))))))))
                     (Disj (LessEqual (Var (Suc (Suc (Suc (Suc (Suc 0)))))) (Var (Suc (Suc (Suc (Suc (Suc 0))))))))
                           (Add (Var (Suc (Suc (Suc (Suc 0)))))))
                           (Minus (Mult (Mult (Var 0) (Inverse (Var (Suc 0))))))
                           (Var (Suc (Suc (Suc (Suc (Suc 0)))))))))))
              (Less (Mult (Power (Add (Var (Suc (Suc (Suc (Suc (Suc 0)))))) 2)
                           (Inverse (Var (Suc (Suc (Suc 0))))))))
                     (Mult (Add (Var (Suc (Suc (Suc 0)))))) (Minus (Var (Suc (Suc 0)))))) (Add (Var (Suc 0)) (Minus (Var 0)))))))
              ([ae, ao, se, so, Interval' (Float 2 0) (Float 2 0), ve, vo, Interval' (Float 0 1) (Float 0 1)] (replicate 8 0))

```

lemma less-Suc-iff-disj: $i < \text{Suc } x \longleftrightarrow i = x \vee i < x$
⟨proof⟩

lemma checker'-soundness-correctness:

```

assumes a ∈ {real-of-float al .. real-of-float au}
assumes b ∈ {real-of-float bl .. real-of-float bu}
assumes c ∈ {real-of-float cl .. real-of-float cu}
assumes d ∈ {real-of-float dl .. real-of-float du}
assumes e ∈ {real-of-float el .. real-of-float eu}
assumes f ∈ {real-of-float fl .. real-of-float fu}
assumes chk: checker' p (Interval' al au) (Interval' bl bu) (Interval' cl cu)
          (Interval' dl du) (Interval' el eu) (Interval' fl fu)
shows checker a b c d e f
⟨proof⟩

```

lemma approximate-soundness-correctness:

```

assumes a ∈ {real-of-float al .. real-of-float au}
assumes b ∈ {real-of-float bl .. real-of-float bu}
assumes c ∈ {real-of-float cl .. real-of-float cu}
assumes d ∈ {real-of-float dl .. real-of-float du}
assumes e ∈ {real-of-float el .. real-of-float eu}
assumes f ∈ {real-of-float fl .. real-of-float fu}
assumes chk: checker' p (Interval' al au) (Interval' bl bu) (Interval' cl cu)
          (Interval' dl du) (Interval' el eu) (Interval' fl fu)
shows checker'-precond: check-precond a b c d e f

```

and *checker'-no-collision: safe-distance.no-collision c b a f e d {0..}*
{proof}

1.5.3 Symbolic Checker

definition *symbolic-checker :: real ⇒ real ⇒ real ⇒ real ⇒ real ⇒ bool*
where

symbolic-checker s_e v_e a_e s_o v_o a_o ≡
let e-stop = - v_e / a_e;
o-stop = - v_o / a_o
in check-precond s_e v_e a_e s_o v_o a_o ∧
 $(\neg \text{quadroot-in } 0 (\min e\text{-stop } o\text{-stop}) (1/2 * (a_o - a_e)) (v_o - v_e) (s_o - s_e)) \wedge$
 $\neg \text{quadroot-in } e\text{-stop } o\text{-stop} (1/2 * a_o) v_o (s_o - \text{movement.p } a_e v_e s_e e\text{-stop})$
∧
 $\neg \text{quadroot-in } o\text{-stop } e\text{-stop} (1/2 * a_e) v_e (s_e - \text{movement.p } a_o v_o s_o o\text{-stop}))$

theorem *symbolic-soundness-correctness:*

symbolic-checker s_e v_e a_e s_o v_o a_o ↔ check-precond s_e v_e a_e s_o v_o a_o ∧
safe-distance.no-collision a_e v_e s_e a_o v_o s_o {0..}}
{proof}
end

end

2 Safe Distance with Reaction Time

theory *Safe-Distance-Reaction*
imports
Safe-Distance
begin

2.1 Normal Safe Distance

locale *safe-distance-normal = safe-distance +*
fixes $\delta :: \text{real}$
assumes *pos-react: $0 < \delta$*
begin

sublocale *ego2: braking-movement a_e v_e (ego.q δ) {proof}*

lemma *ego2-s-init: ego2.s 0 = ego.q δ {proof}*

definition $\tau :: \text{real} \Rightarrow \text{real}$ **where**
 $\tau t = t - \delta$

definition $\tau' :: \text{real} \Rightarrow \text{real}$ **where**
 $\tau' t = 1$

lemma *τ -continuous[continuous-intros]: continuous-on T τ*

$\langle proof \rangle$

lemma *isCont- τ [continuous-intros]*: *isCont τ x*
 $\langle proof \rangle$

lemma *del-has-vector-derivative[derivative-intros]*: (τ has-vector-derivative $\tau' t$)
(at t within u)
 $\langle proof \rangle$

lemma *del-has-real-derivative[derivative-intros]*: (τ has-real-derivative $\tau' t$) (at t within u)
 $\langle proof \rangle$

lemma *delay-image*: $\tau` \{\delta..\} = \{0..\}$
 $\langle proof \rangle$

lemma *s-delayed-has-real-derivative[derivative-intros]*:
assumes $\delta \leq t$
shows $((ego2.s \circ \tau) \text{ has-field-derivative } ego2.s'(t - \delta) * \tau' t)$ (at t within $\{\delta..\}$)
 $\langle proof \rangle$

lemma *s-delayed-has-real-derivative'* [*derivative-intros*]:
assumes $\delta \leq t$
shows $((ego2.s \circ \tau) \text{ has-field-derivative } (ego2.s' \circ \tau) t)$ (at t within $\{\delta..\}$)
 $\langle proof \rangle$

lemma *s-delayed-has-vector-derivative'* [*derivative-intros*]:
assumes $\delta \leq t$
shows $((ego2.s \circ \tau) \text{ has-vector-derivative } (ego2.s' \circ \tau) t)$ (at t within $\{\delta..\}$)
 $\langle proof \rangle$

definition $u :: real \Rightarrow real$ **where**
 $u t = (\begin{array}{ll} \text{if } t \leq 0 \text{ then } s_e \\ \text{else if } t \leq \delta \text{ then } ego.q t \\ \text{else } (ego2.s \circ \tau) t \end{array})$

lemma *init-u*: $t \leq 0 \implies u t = s_e$ $\langle proof \rangle$

lemma *u-delta*: $u \delta = ego2.s 0$
 $\langle proof \rangle$

lemma *q-delta*: $ego.q \delta = ego2.s 0$ $\langle proof \rangle$

definition $u' :: real \Rightarrow real$ **where**
 $u' t = (\text{if } t \leq \delta \text{ then } ego.q' t \text{ else } ego2.s'(t - \delta))$

lemma *u'-delta*: $u' \delta = ego2.s' 0$
 $\langle proof \rangle$

```

lemma q'-delta:  $ego.q' \delta = ego2.s' 0$   $\langle proof \rangle$ 

lemma u-has-real-derivative[derivative-intros]:
  assumes nonneg-t:  $t \geq 0$ 
  shows ( $u$  has-real-derivative  $u' t$ ) (at  $t$  within  $\{0..\}$ )
   $\langle proof \rangle$ 

definition t-stop :: real where
  t-stop =  $ego2.t\text{-stop} + \delta$ 

lemma t-stop-nonneg:  $0 \leq t\text{-stop}$ 
   $\langle proof \rangle$ 

lemma t-stop-pos:  $0 < t\text{-stop}$ 
   $\langle proof \rangle$ 

lemma t-stop-zero:
  assumes t-stop  $\leq x$ 
  assumes  $x \leq \delta$ 
  shows  $v_e = 0$ 
   $\langle proof \rangle$ 

lemma u'-stop-zero:  $u' t\text{-stop} = 0$ 
   $\langle proof \rangle$ 

definition u-max :: real where
  u-max =  $u (ego2.t\text{-stop} + \delta)$ 

lemma u-max-eq:  $u\text{-max} = ego.q \delta - v_e^2 / a_e / 2$ 
   $\langle proof \rangle$ 

lemma u-mono:
  assumes  $x \leq y$   $y \leq t\text{-stop}$ 
  shows  $u x \leq u y$ 
   $\langle proof \rangle$ 

lemma u-antimono:  $x \leq y \implies t\text{-stop} \leq x \implies u y \leq u x$ 
   $\langle proof \rangle$ 

lemma u-max:  $u x \leq u\text{-max}$ 
   $\langle proof \rangle$ 

lemma u-eq-u-stop: NO-MATCH  $t\text{-stop} x \implies x \geq t\text{-stop} \implies u x = u\text{-max}$ 
   $\langle proof \rangle$ 

lemma at-least-delta:
  assumes  $x \leq \delta$ 
  assumes  $t\text{-stop} \leq x$ 
  shows  $ego.q x = ego2.s (x - \delta)$ 

```

$\langle proof \rangle$

lemma *continuous-on-u[continuous-intros]*: *continuous-on T u*
 $\langle proof \rangle$

lemma *isCont-u[continuous-intros]*: *isCont u x*
 $\langle proof \rangle$

definition *collision-react* :: *real set* \Rightarrow *bool* **where**
collision-react time-set \equiv $(\exists t \in \text{time-set}. u t = \text{other.s } t)$

abbreviation *no-collision-react* :: *real set* \Rightarrow *bool* **where**
no-collision-react time-set \equiv $\neg \text{collision-react time-set}$

lemma *no-collision-reactI*:
assumes $\bigwedge t. t \in S \implies u t \neq \text{other.s } t$
shows *no-collision-react S*
 $\langle proof \rangle$

lemma *no-collision-union*:
assumes *no-collision-react S*
assumes *no-collision-react T*
shows *no-collision-react (S \cup T)*
 $\langle proof \rangle$

lemma *collision-trim-subset*:
assumes *collision-react S*
assumes *no-collision-react T*
assumes *T \subseteq S*
shows *collision-react (S - T)*
 $\langle proof \rangle$

theorem *cond-1r* : *u-max < s_o* \implies *no-collision-react {0..}*
 $\langle proof \rangle$

definition *safe-distance-1r* :: *real* **where**
 $\text{safe-distance-1r} = v_e * \delta - v_e^2 / a_e / 2$

lemma *sd-1r-eq*: $(s_o - s_e > \text{safe-distance-1r}) = (u\text{-max} < s_o)$
 $\langle proof \rangle$

lemma *sd-1r-correct*:
assumes $s_o - s_e > \text{safe-distance-1r}$
shows *no-collision-react {0..}*
 $\langle proof \rangle$

lemma *u-other-strict-ivt*:
assumes $u t > \text{other.s } t$
shows *collision-react {0..<t}*

$\langle proof \rangle$

lemma *collision-react-subset*: *collision-react* $s \implies s \subseteq t \implies \text{collision-react } t$
 $\langle proof \rangle$

lemma *u-other-intv*:
 assumes $u \ t \geq \text{other}.s \ t$
 shows *collision-react* $\{0 .. t\}$
 $\langle proof \rangle$

theorem *cond-2r*:
 assumes $u\text{-max} \geq \text{other}.s\text{-stop}$
 shows *collision-react* $\{0 ..\}$
 $\langle proof \rangle$

definition *ego-other2* :: *real* \Rightarrow *real* **where**
 ego-other2 $t = \text{other}.s \ t - u \ t$

lemma *continuous-on-ego-other2*[*continuous-intros*]: *continuous-on* T *ego-other2*
 $\langle proof \rangle$

lemma *isCont-ego-other2*[*continuous-intros*]: *isCont* *ego-other2* x
 $\langle proof \rangle$

definition *ego-other2'* :: *real* \Rightarrow *real* **where**
 ego-other2' $t = \text{other}.s' \ t - u' \ t$

lemma *ego-other2-has-real-derivative*[*derivative-intros*]:
 assumes $0 \leq t$
 shows (*ego-other2 has-real-derivative ego-other2'* t) (at t within $\{0..\}$)
 $\langle proof \rangle$

theorem *cond-3r-1*:
 assumes $u \ \delta \geq \text{other}.s \ \delta$
 shows *collision-react* $\{0 .. \delta\}$
 $\langle proof \rangle$

definition *distance0* :: *real* **where**
 distance0 $= v_e * \delta - v_o * \delta - a_o * \delta^2 / 2$

definition *distance0-2* :: *real* **where**
 distance0-2 $= v_e * \delta + 1 / 2 * v_o^2 / a_o$

theorem *cond-3r-1'*:
 assumes $s_o - s_e \leq \text{distance0}$
 assumes $\delta \leq \text{other}.t\text{-stop}$
 shows *collision-react* $\{0 .. \delta\}$
 $\langle proof \rangle$

```

theorem distance0-2-eq:
  assumes  $\delta > \text{other.t-stop}$ 
  shows  $(u \delta < \text{other.s } \delta) = (s_o - s_e > \text{distance0-2})$ 
  ⟨proof⟩

theorem cond-3r-1'-2:
  assumes  $s_o - s_e \leq \text{distance0-2}$ 
  assumes  $\delta > \text{other.t-stop}$ 
  shows collision-react {0 .. δ}
  ⟨proof⟩

definition safe-distance-3r :: real where
  safe-distance-3r =  $v_e * \delta - v_e^2 / 2 / a_e - v_o * \delta - 1/2 * a_o * \delta^2$ 

lemma distance0-at-most-sd3r:
   $\text{distance0} \leq \text{safe-distance-3r}$ 
  ⟨proof⟩

definition safe-distance-4r :: real where
  safe-distance-4r =  $(v_o + a_o * \delta - v_e)^2 / 2 / (a_o - a_e) - v_o * \delta - 1/2 * a_o * \delta^2 + v_e * \delta$ 

lemma distance0-at-most-sd4r:
  assumes  $a_o > a_e$ 
  shows  $\text{distance0} \leq \text{safe-distance-4r}$ 
  ⟨proof⟩

definition safe-distance-2r :: real where
  safe-distance-2r =  $v_e * \delta - v_e^2 / 2 / a_e + v_o^2 / 2 / a_o$ 

lemma vo-start-geq-ve:
  assumes  $\delta \leq \text{other.t-stop}$ 
  assumes  $\text{other.s' } \delta \geq v_e$ 
  shows  $u \delta < \text{other.s } \delta$ 
  ⟨proof⟩

theorem so-star-stop-leq-se-stop:
  assumes  $\delta \leq \text{other.t-stop}$ 
  assumes  $\text{other.s' } \delta < v_e$ 
  assumes  $\neg (a_o > a_e \wedge \text{other.s' } \delta < v_e \wedge v_e - a_e / a_o * \text{other.s' } \delta < 0)$ 
  shows  $0 \leq -v_e^2 / a_e / 2 + (v_o + a_o * \delta)^2 / a_o / 2$ 
  ⟨proof⟩

theorem distance0-at-most-distance2r:
  assumes  $\delta \leq \text{other.t-stop}$ 
  assumes  $\text{other.s' } \delta < v_e$ 
  assumes  $\neg (a_o > a_e \wedge \text{other.s' } \delta < v_e \wedge v_e - a_e / a_o * \text{other.s' } \delta < 0)$ 
  shows  $\text{distance0} \leq \text{safe-distance-2r}$ 
  ⟨proof⟩

```

```

theorem dist0-sd2r-1:
  assumes  $\delta \leq \text{other.t-stop}$ 
  assumes  $\neg (a_o > a_e \wedge \text{other.s}' \delta < v_e \wedge v_e - a_e / a_o * \text{other.s}' \delta < 0)$ 
  assumes  $s_o - s_e > \text{safe-distance-}2r$ 
  shows  $s_o - s_e > \text{distance0}$ 
  ⟨proof⟩

theorem sd2r-eq:
  assumes  $\delta > \text{other.t-stop}$ 
  shows  $(u\text{-max} < \text{other.s } \delta) = (s_o - s_e > \text{safe-distance-}2r)$ 
  ⟨proof⟩

theorem dist0-sd2r-2:
  assumes  $\delta > -v_o / a_o$ 
  assumes  $s_o - s_e > \text{safe-distance-}2r$ 
  shows  $s_o - s_e > \text{distance0-}2$ 
  ⟨proof⟩
end

```

2.2 Safe Distance Delta

```

locale safe-distance-no-collision-delta = safe-distance-normal +
  assumes no-collision-delta:  $u \delta < \text{other.s } \delta$ 
begin

sublocale delayed-safe-distance: safe-distance  $a_e \ v_e \ \text{ego.q } \delta \ a_o \ \text{other.s}' \delta \ \text{other.s } \delta$ 
  ⟨proof⟩

lemma no-collision-react-initially-strict:
  assumes  $s_o \leq u\text{-max}$ 
  assumes  $u\text{-max} < \text{other.s-stop}$ 
  shows no-collision-react { $0 <..< \delta$ }
  ⟨proof⟩

lemma no-collision-react-initially:
  assumes  $s_o \leq u\text{-max}$ 
  assumes  $u\text{-max} < \text{other.s-stop}$ 
  shows no-collision-react { $0 .. \delta$ }
  ⟨proof⟩

lemma collision-after-delta:
  assumes  $s_o \leq u\text{-max}$ 
  assumes  $u\text{-max} < \text{other.s-stop}$ 
  shows collision-react { $0 .. \delta$ }  $\longleftrightarrow$  collision-react { $\delta .. \delta$ }
  ⟨proof⟩

lemma collision-react-strict:
  assumes  $s_o \leq u\text{-max}$ 

```

```

assumes  $u\text{-max} < \text{other}.s\text{-stop}$ 
shows collision-react  $\{\delta ..\} \longleftrightarrow \text{collision-react } \{\delta <..\}$ 
⟨proof⟩

lemma delayed-other-s-stop-eq: delayed-safe-distance.other.s-stop = other.s-stop
⟨proof⟩

lemma delayed-cond3':
assumes other.s  $\delta \leq u\text{-max}$ 
assumes  $u\text{-max} < \text{other}.s\text{-stop}$ 
shows delayed-safe-distance.collision  $\{0 ..\} \longleftrightarrow$ 
 $(a_o > a_e \wedge \text{other}.s' \delta < v_e \wedge \text{other}.s \delta - \text{ego}.q \delta \leq \text{delayed-safe-distance.snd-safe-distance}$ 
 $\wedge v_e - a_e / a_o * \text{other}.s' \delta < 0)$ 
⟨proof⟩

lemma delayed-other-t-stop-eq:
assumes  $\delta \leq \text{other}.t\text{-stop}$ 
shows delayed-safe-distance.other.t-stop +  $\delta = \text{other}.t\text{-stop}$ 
⟨proof⟩

lemma delayed-other-s-eq:
assumes  $0 \leq t$ 
shows delayed-safe-distance.other.s  $t = \text{other}.s (t + \delta)$ 
⟨proof⟩

lemma translate-collision-range:
assumes  $s_o \leq u\text{-max}$ 
assumes  $u\text{-max} < \text{other}.s\text{-stop}$ 
shows delayed-safe-distance.collision  $\{0 ..\} \longleftrightarrow \text{collision-react } \{\delta ..\}$ 
⟨proof⟩

theorem cond-3r-2:
assumes  $s_o \leq u\text{-max}$ 
assumes  $u\text{-max} < \text{other}.s\text{-stop}$ 
assumes other.s  $\delta \leq u\text{-max}$ 
shows collision-react  $\{0 ..\} \longleftrightarrow$ 
 $(a_o > a_e \wedge \text{other}.s' \delta < v_e \wedge \text{other}.s \delta - \text{ego}.q \delta \leq \text{delayed-safe-distance.snd-safe-distance}$ 
 $\wedge v_e - a_e / a_o * \text{other}.s' \delta < 0)$ 
⟨proof⟩

lemma sd-2r-correct-for-3r-2:
assumes  $s_o - s_e > \text{safe-distance-2r}$ 
assumes other.s  $\delta \leq u\text{-max}$ 
assumes  $\neg (a_o > a_e \wedge \text{other}.s' \delta < v_e \wedge v_e - a_e / a_o * \text{other}.s' \delta < 0)$ 
shows no-collision-react  $\{0..\}$ 
⟨proof⟩

lemma sd2-at-most-sd4:
assumes  $a_o > a_e$ 

```

shows $\text{safe-distance-}2r \leq \text{safe-distance-}4r$
 $\langle proof \rangle$

lemma $sd\text{-}4r\text{-correct}$:

assumes $s_o - s_e > \text{safe-distance-}4r$
assumes $\text{other.s } \delta \leq u\text{-max}$
assumes $\delta \leq \text{other.t-stop}$
assumes $a_o > a_e$
shows $\text{no-collision-react } \{0..\}$

$\langle proof \rangle$

Irrelevant, this Safe Distance is unreachable in the checker.

definition $\text{safe-distance-}5r :: \text{real where}$

$$\text{safe-distance-}5r = v_e^2 / 2 / (a_o - a_e) + v_o^2 / 2 / a_o + v_e * \delta$$

lemma $sd\text{-}5r\text{-correct}$:

assumes $s_o - s_e > \text{safe-distance-}5r$
assumes $u\text{-max} < \text{other.s-stop}$
assumes $\text{other.s } \delta \leq u\text{-max}$
assumes $\delta > \text{other.t-stop}$
shows $\text{no-collision-react } \{0..\}$

$\langle proof \rangle$

lemma $\text{translate-no-collision-range}$:

$$\text{delayed-safe-distance.no-collision } \{0..\} \longleftrightarrow \text{no-collision-react } \{\delta..\}$$

$\langle proof \rangle$

lemma delayed-cond1 :

assumes $\text{other.s } \delta > u\text{-max}$
shows $\text{delayed-safe-distance.no-collision } \{0..\}$

$\langle proof \rangle$

theorem $cond\text{-}3r\text{-}3$:

assumes $s_o \leq u\text{-max}$
assumes $u\text{-max} < \text{other.s-stop}$
assumes $\text{other.s } \delta > u\text{-max}$
shows $\text{no-collision-react } \{0..\}$

$\langle proof \rangle$

lemma $sd\text{-}2r\text{-correct-for-}3r\text{-}3$:

assumes $s_o - s_e > \text{safe-distance-}2r$
assumes $\text{other.s } \delta > u\text{-max}$
shows $\text{no-collision-react } \{0..\}$

$\langle proof \rangle$

lemma $sd\text{-}3r\text{-correct}$:

assumes $s_o - s_e > \text{safe-distance-}3r$
assumes $\delta \leq \text{other.t-stop}$
shows $\text{no-collision-react } \{0..\}$

$\langle proof \rangle$

```

lemma sd-2-at-least-sd-3:
  assumes  $\delta \leq other.t-stop$ 
  shows safe-distance-3r  $\geq$  safe-distance-2r
⟨proof⟩
end

```

2.3 Checker Design

We define two checkers for different cases:

- one checker for the case that $\delta \leq other.t-stop$ ($other.t-stop = - v_o / a_o$)
- a second checker for the case that $\delta > other.t-stop$

```

definition check-precond-r1 :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real
 $\Rightarrow$  bool where
  check-precond-r1  $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta \longleftrightarrow s_o > s_e \wedge 0 \leq v_e \wedge 0 \leq v_o \wedge a_e < 0$ 
 $\wedge a_o < 0 \wedge 0 < \delta \wedge \delta \leq - v_o / a_o$ 

```

```

definition safe-distance0 :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real where
  safe-distance0  $v_e \ a_o \ v_o \ \delta = v_e * \delta - v_o * \delta - a_o * \delta^2 / 2$ 

```

```

definition safe-distance-1r :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real where
  safe-distance-1r  $a_e \ v_e \ \delta = v_e * \delta - v_e^2 / a_e / 2$ 

```

```

definition safe-distance-2r :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real where
  safe-distance-2r  $a_e \ v_e \ a_o \ v_o \ \delta = v_e * \delta - v_e^2 / 2 / a_e + v_o^2 / 2 / a_o$ 

```

```

definition safe-distance-4r :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real where
  safe-distance-4r  $a_e \ v_e \ a_o \ v_o \ \delta = (v_o + a_o * \delta - v_e)^2 / 2 / (a_o - a_e) - v_o * \delta$ 
 $- 1 / 2 * a_o * \delta^2 + v_e * \delta$ 

```

```

definition safe-distance-3r :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real where
  safe-distance-3r  $a_e \ v_e \ a_o \ v_o \ \delta = v_e * \delta - v_e^2 / 2 / a_e - v_o * \delta - 1 / 2 * a_o * \delta^2$ 

```

```

definition checker-r1 :: real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  real  $\Rightarrow$  bool
where
  checker-r1  $s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta \equiv$ 
    let distance  $= s_o - s_e;$ 
    precond  $= check-precond-r1 \ s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta;$ 
    vo-star  $= v_o + a_o * \delta;$ 
    t-stop-o-star  $= - vo-star / a_o;$ 
    t-stop-e  $= - v_e / a_e;$ 
    safe-dist0  $= safe-distance-1r \ a_e \ v_e \ \delta;$ 
    safe-dist1  $= safe-distance-2r \ a_e \ v_e \ a_o \ v_o \ \delta;$ 

```

```

safe-dist2   = safe-distance-4r ae ve ao vo δ;
safe-dist3   = safe-distance-3r ae ve ao vo δ
in
  if  $\neg$  precond then False
  else if distance > safe-dist0  $\vee$  distance > safe-dist3 then True
  else if (ao > ae  $\wedge$  vo-star < ve  $\wedge$  t-stop-e < t-stop-o-star) then distance >
safe-dist2
  else distance > safe-dist1

```

theorem checker-r1-correctness:

```
(checker-r1 se ve ae so vo ao δ  $\longleftrightarrow$  check-precond-r1 se ve ae so vo ao δ  $\wedge$ 
safe-distance-normal.no-collision-react ae ve se ao vo so δ {0..})
```

(proof)

definition check-precond-r2 :: real \Rightarrow bool **where**

$$\text{check-precond-r2 } s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta \longleftrightarrow s_o > s_e \ \wedge \ 0 \leq v_e \ \wedge \ 0 \leq v_o \ \wedge \ a_e < 0 \ \wedge \ a_o < 0 \ \wedge \ 0 < \delta \ \wedge \ \delta > -v_o / a_o$$

definition safe-distance0-2 :: real \Rightarrow real \Rightarrow real \Rightarrow real **where**

$$\text{safe-distance0-2 } v_e \ a_o \ v_o \ \delta = v_e * \delta + 1 / 2 * v_o^2 / a_o$$

definition checker-r2 :: real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow real \Rightarrow bool **where**

$$\begin{aligned} \text{checker-r2 } s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta &\equiv \\ \text{let } distance &= s_o - s_e; \\ \text{precond} &= \text{check-precond-r2 } s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta; \\ \text{safe-dist0} &= \text{safe-distance-1r } a_e \ v_e \ \delta; \\ \text{safe-dist1} &= \text{safe-distance-2r } a_e \ v_e \ a_o \ v_o \ \delta \end{aligned}$$

in
 if \neg precond then False
 else if distance > safe-dist0 then True
 else distance > safe-dist1

theorem checker-r2-correctness:

```
(checker-r2 se ve ae so vo ao δ  $\longleftrightarrow$  check-precond-r2 se ve ae so vo ao δ  $\wedge$ 
safe-distance-normal.no-collision-react ae ve se ao vo so δ {0..})
```

(proof)

Combine the two checkers into one.

definition check-precond-r :: real \Rightarrow bool **where**

$$\text{check-precond-r } s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta \longleftrightarrow s_o > s_e \ \wedge \ 0 \leq v_e \ \wedge \ 0 \leq v_o \ \wedge \ a_e < 0 \ \wedge \ a_o < 0 \ \wedge \ 0 < \delta$$

definition checker-r :: real \Rightarrow bool **where**

$$\begin{aligned} \text{checker-r } s_e \ v_e \ a_e \ s_o \ v_o \ a_o \ \delta &\equiv \\ \text{let } distance &= s_o - s_e; \end{aligned}$$

```

precond      = check-precond-r se ve ae so vo ao δ;
vo-star     = vo + ao * δ;
t-stop-o-star = -vo-star / ao;
t-stop-e     = -ve / ae;
t-stop-o     = -vo / ao;
safe-dist0   = safe-distance-1r ae ve δ;
safe-dist1   = safe-distance-2r ae ve ao vo δ;
safe-dist2   = safe-distance-4r ae ve ao vo δ;
safe-dist3   = safe-distance-3r ae ve ao vo δ
in
if  $\neg$  precond then False
else if distance > safe-dist0 then True
else if  $\delta \leq t\text{-stop-}o \wedge distance > safe\text{-}dist3$  then True
else if  $\delta \leq t\text{-stop-}o \wedge (a_o > a_e \wedge vo\text{-}star < v_e \wedge t\text{-stop-}e < t\text{-stop-}o\text{-}star)$ 
then distance > safe-dist2
else distance > safe-dist1

```

theorem checker-eq-1:

```

checker-r se ve ae so vo ao δ  $\wedge \delta \leq -v_o / a_o \longleftrightarrow$  checker-r1 se ve ae so vo ao
δ
⟨proof⟩

```

theorem checker-eq-2:

```

checker-r se ve ae so vo ao δ  $\wedge \delta > -v_o / a_o \longleftrightarrow$  checker-r2 se ve ae so vo ao δ
⟨proof⟩

```

theorem checker-r-correctness:

```

(checker-r se ve ae so vo ao δ  $\longleftrightarrow$  check-precond-r se ve ae so vo ao δ  $\wedge$ 
safe-distance-normal.no-collision-react ae ve se ao vo so δ {0..})
⟨proof⟩

```

end

3 Evaluation

theory Evaluation

imports

Safe-Distance

HOL-Library.Float

begin

3.1 Code Generation Setup for Numeric Values

definition real-div-down :: nat \Rightarrow int \Rightarrow int \Rightarrow real **where**

real-div-down p i j = truncate-down (Suc p) (i / j)

definition real-div-up :: nat \Rightarrow int \Rightarrow int \Rightarrow real **where**

real-div-up p i j = truncate-up (Suc p) (i / j)

```

context includes float.lifting begin
lift-definition float-div-down :: nat  $\Rightarrow$  int  $\Rightarrow$  int  $\Rightarrow$  float is real-div-down
   $\langle proof \rangle$ 

lift-definition float-div-up :: nat  $\Rightarrow$  int  $\Rightarrow$  int  $\Rightarrow$  float is real-div-up
   $\langle proof \rangle$ 
end

lemma compute-float-div-up[code]: float-div-up p i j = - float-div-down p (-i) j
  including float.lifting
   $\langle proof \rangle$ 

lemma compute-float-div-down[code]:
  float-div-down prec m1 m2 = lapprox-rat (Suc prec) m1 m2
  including float.lifting  $\langle proof \rangle$ 

definition real2-of-real :: nat  $\Rightarrow$  real  $\Rightarrow$  (real * real) where
  real2-of-real p x = (truncate-down (Suc p) x, truncate-up (Suc p) x)

context includes float.lifting begin
lift-definition float2-of-real :: nat  $\Rightarrow$  real  $\Rightarrow$  float  $\times$  float is real2-of-real
   $\langle proof \rangle$ 
end

definition float2-opt-of-real :: nat  $\Rightarrow$  real  $\Rightarrow$  float interval option where
  float2-opt-of-real prec x = Interval' (fst (float2-of-real prec x)) (snd (float2-of-real
  prec x))

hide-const (open) Fraction-Field.Fract
lemma real-of-rat-Fract[simp]: real-of-rat (Fract a b) = a / b
   $\langle proof \rangle$ 

lemma [code]: float2-of-real p (Ratreal r) =
  (let (a, b) = quotient-of r in
  (float-div-down p a b, float-div-up p a b))
  including float.lifting
   $\langle proof \rangle$ 

fun real-of-dec :: integer  $\times$  integer  $\Rightarrow$  real where
  real-of-dec (m, e) =
    real-of-int (int-of-integer m) *
      (if e  $\geq$  0 then 10  $^$  (nat-of-integer e) else inverse (10  $^$  (nat (-(int-of-integer
      e))))))

lemma real-of-dec (m, e) = int-of-integer m * 10 powr (int-of-integer e)
   $\langle proof \rangle$ 

```

3.2 Data Evaluation

```
definition trans6 where
  trans6 c chk se ve ae so vo ao =
    chk (c se) (c ve) (c ae) (c so) (c vo) (c ao)

definition checker-dec where
  checker-dec chk p u =
    trans6 (float2-opt-of-real (nat-of-integer u) o real-of-dec) (chk (nat-of-integer
  p))

definition checker-interval = checker-dec checker'
definition checker-symbolic = trans6 real-of-dec symbolic-checker
definition checker-rational = trans6 real-of-dec checker
lemmas[code] = movement.p-def
```

$\langle ML \rangle$

The precision of the input data is roughly 12 and yields similar performance as Sturm

$\langle ML \rangle$

Number of data points:

- data01: 1121215
- data02: 1341135
- data03: 1452656

$\langle ML \rangle$

Precision: 12, Uncertainty: 7 digits

$\langle ML \rangle$

end

References

- [1] A. Rizaldi, F. Immler, and M. Althoff. A formally verified checker of the safe distance traffic rules for autonomous vehicles. In S. Rayadurgam and O. Tkachuk, editors, *NASA Formal Methods*, pages 175–190, Cham, 2016. Springer International Publishing.