

Providing restriction Spaces with an ultrametric Structure

Benoît Ballenghien

Benjamin Puyobro

Burkhart Wolff

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Abstract

We investigate the relationship between restriction spaces and classical metric structures by instantiating the former as ultrametric spaces. This is classically captured by defining the distance as

$$\text{dist } x \ y = \inf_{x \downarrow n = y \downarrow n} \left(\frac{1}{2} \right)^n$$

but we actually generalize this perspective by introducing a hierarchy of increasingly refined type classes to systematically relate ultrametric and restriction-based notions. This layered approach enables a precise comparison of structural and topological properties. In the end, our main result establishes that completeness in the sense of restriction spaces coincides with standard metric completeness, thus bridging the gap between `Restriction_Spaces` and Banach's fixed-point theorem established in `HOL-Analysis`.

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1 Definitions on Functions of Metric Space

In this theory, we define the notion of lipschitz map, non-expanding map and contraction map. We also establish correspondences.

1.1 Definitions

1.1.1 Lipschitz Map

This notion is a generalization of contraction map and non-expanding map.

definition *lipschitz-with-on* :: $\langle [a :: \text{metric-space} \Rightarrow b :: \text{metric-space}, \text{real}, 'a \text{ set}] \Rightarrow \text{bool} \rangle$
where $\langle \text{lipschitz-with-on } f \alpha A \equiv 0 \leq \alpha \wedge (\forall x \in A. \forall y \in A. \text{dist}(f x) (f y) \leq \alpha * \text{dist } x y) \rangle$

abbreviation *lipschitz-with* :: $\langle [a :: \text{metric-space} \Rightarrow b :: \text{metric-space}, \text{real}] \Rightarrow \text{bool} \rangle$
where $\langle \text{lipschitz-with } f \alpha \equiv \text{lipschitz-with-on } f \alpha \text{ UNIV} \rangle$

lemma *lipschitz-with-onI* :
 $\langle [\exists 0 \leq \alpha; \forall x y. [x \in A; y \in A; x \neq y; f x \neq f y] \implies \text{dist}(f x) (f y) \leq \alpha * \text{dist } x y] \implies$
 $\text{lipschitz-with-on } f \alpha A$
 $\langle \text{proof} \rangle$

lemma *lipschitz-withI* :
 $\langle [\exists 0 \leq \alpha; \forall x y. x \neq y \implies f x \neq f y \implies \text{dist}(f x) (f y) \leq \alpha * \text{dist } x y] \implies$
 $\text{lipschitz-with } f \alpha$
 $\langle \text{proof} \rangle$

```

lemma lipschitz-with-onD1 : <lipschitz-with-on f α A => 0 ≤ α>
⟨proof⟩

lemma lipschitz-withD1 : <lipschitz-with f α => 0 ≤ α>
⟨proof⟩

lemma lipschitz-with-onD2 :
<lipschitz-with-on f α A => x ∈ A => y ∈ A => dist (f x) (f y) ≤
α * dist x y>
⟨proof⟩

lemma lipschitz-withD2 :
<lipschitz-with f α => dist (f x) (f y) ≤ α * dist x y>
⟨proof⟩

lemma lipschitz-with-imp-lipschitz-with-on: <lipschitz-with f α => lip-
schitz-with-on f α A>
⟨proof⟩

lemma lipschitz-with-on-imp-lipschitz-with-on-ge : <lipschitz-with-on f
β A>
if ⟨α ≤ β⟩ and <lipschitz-with-on f α A>
⟨proof⟩

theorem lipschitz-with-on-comp-lipschitz-with-on :
<lipschitz-with-on (λx. g (f x)) (β * α) A>
if ⟨f ‘ A ⊆ B⟩ <lipschitz-with-on g β B> <lipschitz-with-on f α A>
⟨proof⟩

corollary lipschitz-with-comp-lipschitz-with :
<[lipschitz-with g β; lipschitz-with f α] =>
lipschitz-with (λx. g (f x)) (β * α)>
⟨proof⟩

1.1.2 Non-expanding Map

definition non-expanding-on :: <['a :: metric-space => 'b :: metric-space,
'a set] => bool>
where <non-expanding-on f A ≡ lipschitz-with-on f 1 A>

abbreviation non-expanding :: <['a :: metric-space => 'b :: metric-space]
⇒ bool>
where <non-expanding f ≡ non-expanding-on f UNIV>

lemma non-expanding-onI :
<[Λx y. [x ∈ A; y ∈ A; x ≠ y; f x ≠ f y] => dist (f x) (f y) ≤ dist

```

```

 $x \ y] \implies$ 
  non-expanding-on  $f A$ 
  ⟨proof⟩

lemma non-expandingI :
  ⟨ $\lambda x \ y. \ x \neq y \implies f x \neq f y \implies \text{dist} (f x) (f y) \leq \text{dist} x y$ ⟩  $\implies$ 
  non-expanding  $f$ 
  ⟨proof⟩

lemma non-expanding-onD :
  ⟨non-expanding-on  $f A \implies x \in A \implies y \in A \implies \text{dist} (f x) (f y) \leq$ 
   $\text{dist} x y$ ⟩
  ⟨proof⟩

lemma non-expandingD : ⟨non-expanding  $f \implies \text{dist} (f x) (f y) \leq \text{dist}$ 
 $x y$ ⟩
  ⟨proof⟩

lemma non-expanding-imp-non-expanding-on: ⟨non-expanding  $f \implies$ 
non-expanding-on  $f A$ ⟩
  ⟨proof⟩

lemma non-expanding-on-comp-non-expanding-on :
  ⟨ $f : A \subseteq B; \text{non-expanding-on } g B; \text{non-expanding-on } f A$ ⟩  $\implies$ 
  non-expanding-on  $(\lambda x. \ g (f x)) A$ 
  ⟨proof⟩

corollary non-expanding-comp-non-expanding :
  ⟨ $\text{non-expanding } g; \text{non-expanding } f$ ⟩  $\implies$  non-expanding  $(\lambda x. \ g (f x))$ 
  ⟨proof⟩

1.1.3 Contraction Map

definition contraction-with-on :: ⟨ $'a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space, real, } 'a \text{ set}$ ⟩  $\Rightarrow$  bool
  where ⟨ $\text{contraction-with-on } f \alpha A \equiv \alpha < 1 \wedge \text{lipschitz-with-on } f \alpha A$ ⟩

abbreviation contraction-with :: ⟨ $'a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space, real}$ ⟩  $\Rightarrow$  bool
  where ⟨ $\text{contraction-with } f \alpha \equiv \text{contraction-with-on } f \alpha \text{ UNIV}$ ⟩

definition contraction-on :: ⟨ $'a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space, }$ 
 $'a \text{ set}$ ⟩  $\Rightarrow$  bool
  where ⟨ $\text{contraction-on } f A \equiv \exists \alpha. \text{contraction-with-on } f \alpha A$ ⟩

abbreviation contraction :: ⟨ $'a :: \text{metric-space} \Rightarrow 'b :: \text{metric-space}$ ⟩

```

```

 $\Rightarrow \text{bool}$ 
where  $\langle \text{contraction } f \equiv \text{contraction-on } f \text{ UNIV} \rangle$ 

lemma  $\text{contraction-with-onI} :$ 
 $\langle \llbracket 0 \leq \alpha; \alpha < 1; \bigwedge x y. \llbracket x \in A; y \in A; x \neq y; f x \neq f y \rrbracket \implies \text{dist}(f x)(f y) \leq \alpha * \text{dist}(x)(y) \rrbracket \implies \text{contraction-with-on } f \alpha A$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-withI} :$ 
 $\langle \llbracket 0 \leq \alpha; \alpha < 1; \bigwedge x y. x \neq y \implies f x \neq f y \implies \text{dist}(f x)(f y) \leq \alpha * \text{dist}(x)(y) \rrbracket \implies \text{contraction-with } f \alpha$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-with-onD1} : \langle \text{contraction-with-on } f \alpha A \implies 0 \leq \alpha \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-withD1} : \langle \text{contraction-with } f \alpha \implies 0 \leq \alpha \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-with-onD2} : \langle \text{contraction-with-on } f \alpha A \implies \alpha < 1 \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-withD2} : \langle \text{contraction-with } f \alpha \implies \alpha < 1 \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-with-onD3} :$ 
 $\langle \text{contraction-with-on } f \alpha A \implies x \in A \implies y \in A \implies \text{dist}(f x)(f y) \leq \alpha * \text{dist}(x)(y) \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-withD3} : \langle \text{contraction-with } f \alpha \implies \text{dist}(f x)(f y) \leq \alpha * \text{dist}(x)(y) \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-with-imp-contraction-with-on} :$ 
 $\langle \text{contraction-with } f \alpha \implies \text{contraction-with-on } f \alpha A \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-imp-contraction-on} : \langle \text{contraction } f \implies \text{contraction-on } f A \rangle$ 
 $\langle \text{proof} \rangle$ 

lemma  $\text{contraction-with-on-imp-contraction-on} :$ 

```

$\langle \text{contraction-with-on } f \alpha A \implies \text{contraction-on } f A \rangle$
 $\langle \text{proof} \rangle$

lemma *contraction-with-imp-contraction*: $\langle \text{contraction-with } f \alpha \implies \text{contraction } f \rangle$
 $\langle \text{proof} \rangle$

lemma *contraction-onE*:
 $\langle \llbracket \text{contraction-on } f A; \wedge \alpha. \text{contraction-with-on } f \alpha A \implies \text{thesis} \rrbracket \implies \text{thesis} \rangle$
 $\langle \text{proof} \rangle$

lemma *contractionE*:
 $\langle \llbracket \text{contraction } f; \wedge \alpha. \text{contraction-with } f \alpha \implies \text{thesis} \rrbracket \implies \text{thesis} \rangle$
 $\langle \text{proof} \rangle$

lemma *contraction-with-on-imp-contraction-with-on-ge* :
 $\langle \llbracket \alpha \leq \beta; \beta < 1; \text{contraction-with-on } f \alpha A \rrbracket \implies \text{contraction-with-on } f \beta A \rangle$
 $\langle \text{proof} \rangle$

1.2 Properties

lemma *contraction-with-on-imp-lipschitz-with-on[simp]* :
 $\langle \text{contraction-with-on } f \alpha A \implies \text{lipschitz-with-on } f \alpha A \rangle$
 $\langle \text{proof} \rangle$

lemma *non-expanding-on-imp-lipschitz-with-one-on[simp]* :
 $\langle \text{non-expanding-on } f A \implies \text{lipschitz-with-on } f 1 A \rangle$
 $\langle \text{proof} \rangle$

lemma *contraction-on-imp-non-expanding-on[simp]* :
 $\langle \text{contraction-on } f A \implies \text{non-expanding-on } f A \rangle$
 $\langle \text{proof} \rangle$

lemma *contraction-with-on-comp-contraction-with-on* :
 $\langle \text{contraction-with-on } (\lambda x. g(f x)) (\beta * \alpha) A \rangle$
if $\langle f ` A \subseteq B \rangle$ $\langle \text{contraction-with-on } g \beta B \rangle$ $\langle \text{contraction-with-on } f \alpha A \rangle$
 $\langle \text{proof} \rangle$

corollary *contraction-with-comp-contraction-with* :
 $\langle \llbracket \text{contraction-with } g \beta; \text{contraction-with } f \alpha \rrbracket \implies \text{contraction-with } (\lambda x. g(f x)) (\beta * \alpha) \rangle$
 $\langle \text{proof} \rangle$

corollary *contraction-on-comp-contraction-on* :
 $\langle \llbracket f ' A \subseteq B; \text{contraction-on } g B; \text{contraction-on } f A \rrbracket \implies \text{contraction-on } (\lambda x. g (f x)) A \rangle$
 $\langle \text{proof} \rangle$

corollary *contraction-comp-contraction* :
 $\langle \llbracket \text{contraction } g; \text{contraction } f \rrbracket \implies \text{contraction } (\lambda x. g (f x)) \rangle$
 $\langle \text{proof} \rangle$

lemma *contraction-with-on-comp-non-expanding-on* :
 $\langle \text{contraction-with-on } (\lambda x. g (f x)) \beta A \rangle$
if $\langle f ' A \subseteq B \rangle$ $\langle \text{contraction-with-on } g \beta B \rangle$ $\langle \text{non-expanding-on } f A \rangle$
 $\langle \text{proof} \rangle$

corollary *contraction-with-comp-non-expanding* :
 $\langle \llbracket \text{contraction-with } g \beta; \text{non-expanding } f \rrbracket \implies \text{contraction-with } (\lambda x. g (f x)) \beta \rangle$
 $\langle \text{proof} \rangle$

corollary *contraction-on-comp-non-expanding-on* :
 $\langle \llbracket f ' A \subseteq B; \text{contraction-on } g B; \text{non-expanding-on } f A \rrbracket \implies \text{contraction-on } (\lambda x. g (f x)) A \rangle$
 $\langle \text{proof} \rangle$

corollary *contraction-comp-non-expanding* :
 $\langle \llbracket \text{contraction } g; \text{non-expanding } f \rrbracket \implies \text{contraction } (\lambda x. g (f x)) \rangle$
 $\langle \text{proof} \rangle$

lemma *non-expanding-on-comp-contraction-with-on* :
 $\langle \text{contraction-with-on } (\lambda x. g (f x)) \alpha A \rangle$
if $\langle f ' A \subseteq B \rangle$ $\langle \text{non-expanding-on } g B \rangle$ $\langle \text{contraction-with-on } f \alpha A \rangle$
 $\langle \text{proof} \rangle$

corollary *non-expanding-comp-contraction-with* :
 $\langle \llbracket \text{non-expanding } g; \text{contraction-with } f \alpha \rrbracket \implies \text{contraction-with } (\lambda x. g (f x)) \alpha \rangle$
 $\langle \text{proof} \rangle$

corollary *non-expanding-on-comp-contraction-on* :
 $\langle \llbracket f ' A \subseteq B; \text{non-expanding-on } g B; \text{contraction-on } f A \rrbracket \implies \text{contraction-on } (\lambda x. g (f x)) A \rangle$
 $\langle \text{proof} \rangle$

corollary *non-expanding-comp-contraction* :
 $\langle \llbracket \text{non-expanding } g; \text{contraction } f \rrbracket \implies \text{contraction } (\lambda x. g (f x)) \rangle$
 $\langle \text{proof} \rangle$

1.3 Banach's fixed-point Theorems

We rewrite the Banach's fixed-point theorems with our new definition.

theorem *Banach-fix-type* : $\langle \text{contraction } f \implies \exists!x. f x = x \rangle$
for $f :: \langle 'a :: \text{complete-space} \Rightarrow 'a \rangle$
 $\langle \text{proof} \rangle$

theorem *Banach-fix*:
 $\langle \text{contraction-on } f s \implies \exists!x. x \in s \wedge f x = x \rangle$ **if** $\langle \text{complete } s \rangle \langle s \neq \{\} \rangle \langle f ` s \subseteq s \rangle$
 $\langle \text{proof} \rangle$

2 Locales factorizing the proof Work

2.1 Preliminaries on strictly decreasing Sequences

abbreviation *strict-decseq* :: $\langle (\text{nat} \Rightarrow 'a :: \text{order}) \Rightarrow \text{bool} \rangle$
where $\langle \text{strict-decseq} \equiv \text{monotone } (<) \ (\lambda x y. y < x) \rangle$

lemma *strict-decseq-def* : $\langle \text{strict-decseq } X \longleftrightarrow (\forall m n. m < n \longrightarrow X n < X m) \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseqI* : $\langle \text{strict-decseq } X \rangle$ **if** $\langle \bigwedge n. X (\text{Suc } n) < X n \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseqD* : $\langle \text{strict-decseq } X \implies m < n \implies X n < X m \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseq-def-bis* : $\langle \text{strict-decseq } X \longleftrightarrow (\forall m n. X n < X m \longleftrightarrow m < n) \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseq-def-ter* : $\langle \text{strict-decseq } X \longleftrightarrow (\forall m n. X n \leq X m \longleftrightarrow m \leq n) \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseq-imp-decseq* : $\langle \text{strict-decseq } \sigma \implies \text{decseq } \sigma \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseq-SucI* : $\langle (\bigwedge n. X (\text{Suc } n) < X n) \implies \text{strict-decseq } X \rangle$
 $\langle \text{proof} \rangle$

lemma *strict-decseq-SucD* : $\langle \text{strict-decseq } A \implies A (\text{Suc } i) < A i \rangle$
 $\langle \text{proof} \rangle$

Classically, a restriction space is given the structure of a metric space by defining $\text{dist } x \ y \equiv \text{Inf} \{(1 / 2)^n \mid n. \ x \downarrow n = y \downarrow n\}$. This obviously also works if we replace $1 / 2$ by any real δ such that $0 < \delta$ and $\delta < 1$. But more generally, this still works if we set $\text{dist } x \ y \equiv \text{Inf} \{\sigma_n \mid n. \ x \downarrow n = y \downarrow n\}$ where σ is a sequence of *real* verifying $\forall n. \ 0 < \sigma_n$ and $\sigma \longrightarrow 0$. As you would expect, the more structure you have, the more powerful theorems you get. We explore all these variants in the theory below.

2.2 The Construction with Locales

Our formalization will extend the class *metric-space*. But some proofs are redundant, especially when it comes to the product type. So first we will be working with locales, and interpret them with the classes.

```

locale NonDecseqRestrictionSpace = PreorderRestrictionSpace +
  — Factorization of the proof work.
  fixes M :: ‹'a set› and restriction- $\sigma$  :: ‹nat  $\Rightarrow$  real› (⟨ $\sigma_\downarrow$ ⟩)
  and restriction-dist :: ‹'a  $\Rightarrow$  'a  $\Rightarrow$  real› (⟨ $\text{dist}_\downarrow$ ⟩)
  assumes restriction- $\sigma$ -tends-to-zero : ⟨restriction- $\sigma$   $\longrightarrow$  0⟩
  and zero-less-restriction- $\sigma$  [simp] : ⟨0 < restriction- $\sigma$  n⟩
  and dist-restriction-is :
    ⟨ $\text{dist}_\downarrow x \ y = (\text{INF } n \in \text{restriction-related-set } x \ y. \ \text{restriction-}\sigma \ n)$ ⟩
begin

  lemma zero-le-restriction- $\sigma$  [simp] : ⟨0 ≤  $\sigma_\downarrow n$ ⟩
    ⟨proof⟩

  lemma restriction- $\sigma$ -neq-zero [simp] : ⟨ $\sigma_\downarrow n \neq 0$ ⟩
    ⟨proof⟩

  lemma bounded-range-restriction- $\sigma$ : ⟨bounded (range  $\sigma_\downarrow$ )⟩
    ⟨proof⟩

  abbreviation restriction- $\sigma$ -related-set :: ‹'a  $\Rightarrow$  'a  $\Rightarrow$  real set›
    where ⟨restriction- $\sigma$ -related-set x y ≡  $\sigma_\downarrow` \text{restriction-related-set } x \ y$ ⟩

  abbreviation restriction- $\sigma$ -not-related-set :: ‹'a  $\Rightarrow$  'a  $\Rightarrow$  real set›
    where ⟨restriction- $\sigma$ -not-related-set x y ≡  $\sigma_\downarrow` \text{restriction-not-related-set } x \ y$ ⟩

```

```

lemma nonempty-restriction- $\sigma$ -related-set :
  ⟨restriction- $\sigma$ -related-set  $x y \neq \{\}$ ⟩ ⟨proof⟩

lemma restriction- $\sigma$ -related-set-Un-restriction- $\sigma$ -not-related-set :
  ⟨restriction- $\sigma$ -related-set  $x y \cup$  restriction- $\sigma$ -not-related-set  $x y =$ 
  range  $\sigma_\downarrow$ ⟩
  ⟨proof⟩

lemma ⟨bdd-above (restriction- $\sigma$ -related-set  $x y$ )⟩
  ⟨proof⟩

lemma ⟨bdd-above (restriction- $\sigma$ -not-related-set  $x y$ )⟩
  ⟨proof⟩

lemma bounded-restriction- $\sigma$ -related-set: ⟨bounded (restriction- $\sigma$ -related-set
 $x y$ )⟩
  ⟨proof⟩

lemma bounded-restriction- $\sigma$ -not-related-set: ⟨bounded (restriction- $\sigma$ -not-related-set
 $x y$ )⟩
  ⟨proof⟩

corollary restriction-space-Inf-properties:
  ⟨ $a \in$  restriction- $\sigma$ -related-set  $x y \implies dist_\downarrow x y \leq a$ ⟩
  ⟨[ $\forall a. a \in$  restriction- $\sigma$ -related-set  $x y \implies b \leq a$ ]  $\implies b \leq dist_\downarrow x y$ ⟩
  ⟨proof⟩

lemma restriction- $\sigma$ -related-set-alt :
  ⟨restriction- $\sigma$ -related-set  $x y = \{\sigma_\downarrow n | n. n \in$  restriction-related-set
 $x y\}$ ⟩
  ⟨proof⟩

lemma exists-less-restriction- $\sigma$  : ⟨ $\exists n. m < n \wedge \sigma_\downarrow n < \sigma_\downarrow m$ ⟩
  ⟨proof⟩

lemma ⟨ $dist_\downarrow x y = Inf$  (restriction- $\sigma$ -related-set  $x y$ )⟩
  ⟨proof⟩

lemma not-related-imp-dist-restriction-is-some-restriction- $\sigma$  :
  ⟨ $\exists n. dist_\downarrow x y = \sigma_\downarrow n \wedge (\forall m \leq n. x \downarrow m \lessapprox y \downarrow m) \wedge$  if ⟨ $\neg x \lessapprox y$ ⟩
  ⟨proof⟩

lemma not-related-imp-dist-restriction-le-some-restriction- $\sigma$  :
  ⟨ $\neg x \lessapprox y \implies \exists n. dist_\downarrow x y \leq \sigma_\downarrow n \wedge (\neg x \downarrow Suc n \lessapprox y \downarrow Suc n) \wedge$ 
  ( $\forall m \leq n. x \downarrow m \lessapprox y \downarrow m$ )⟩

```

$\langle proof \rangle$

lemma *restriction-dist-eq-0-iff-related* : $\langle dist_{\downarrow} x y = 0 \longleftrightarrow x \lesssim y \rangle$
 $\langle proof \rangle$

end

locale *DecseqRestrictionSpace* = *NonDecseqRestrictionSpace* +
 assumes *decseq-restriction-σ* : $\langle decseq \sigma_{\downarrow} \rangle$
begin

lemma *dist-restriction-is-bis* :
 $\langle dist_{\downarrow} x y = (if x \lesssim y then 0 else \sigma_{\downarrow} (Sup (restriction-related-set x y))) \rangle$
 if $\langle x \in M \rangle$ **and** $\langle y \in M \rangle$
 $\langle proof \rangle$

lemma *not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq* :
 $\langle dist_{\downarrow} x y = \sigma_{\downarrow} (Sup (restriction-related-set x y)) \rangle$
 $\langle \forall m \leq Sup (restriction-related-set x y). x \downarrow m \lesssim y \downarrow m \rangle$
 $\langle \forall m > Sup (restriction-related-set x y). \neg x \downarrow m \lesssim y \downarrow m \rangle$
 if $\langle \neg x \lesssim y \rangle$ **and** $\langle x \in M \rangle$ **and** $\langle y \in M \rangle$
 $\langle proof \rangle$

theorem *restriction-dist-tends-to-zero-independent-of-restriction-σ* :
 — Very powerful theorem: the convergence of the distance to 0 is
 actually independent from the restriction sequence chosen.
 — This is the theorem for which we had to work with locales first.

assumes $\langle DecseqRestrictionSpace (\downarrow) (\lesssim) restriction-\sigma' restriction-dist' \rangle$
 and $\langle \Sigma \in M \rangle$ **and** $\langle range \sigma \subseteq M \rangle$
shows $\langle (\lambda n. dist_{\downarrow} (\sigma n) \Sigma) \longrightarrow 0 \longleftrightarrow (\lambda n. restriction-dist' (\sigma n) \Sigma) \longrightarrow 0 \rangle$
 $\langle proof \rangle$

end

3 Ultrametric Structure of restriction Spaces

This has only be proven with the sort constraint, not inside the context of the class *metric-space* ...

context *metric-space* **begin**

lemma *LIMSEQ-def* : $\langle X \longrightarrow L \longleftrightarrow (\forall r>0. \exists no. \forall n \geq no. dist(X n) L < r) \rangle$
 $\langle proof \rangle$

lemma *LIMSEQ-iff-nz*: $\langle X \longrightarrow L \longleftrightarrow (\forall r>0. \exists no>0. \forall n \geq no. dist(X n) L < r) \rangle$
 $\langle proof \rangle$

lemma *metric-LIMSEQ-I*: $\langle (\bigwedge r. 0 < r \implies \exists no. \forall n \geq no. dist(X n) L < r) \implies X \longrightarrow L \rangle$
 $\langle proof \rangle$

lemma *metric-LIMSEQ-D*: $\langle X \longrightarrow L \implies 0 < r \implies \exists no. \forall n \geq no. dist(X n) L < r \rangle$
 $\langle proof \rangle$

lemma *LIMSEQ-dist-iff*:
 $\langle f \longrightarrow l \longleftrightarrow (\lambda x. dist(f x) l) \longrightarrow 0 \rangle$
 $\langle proof \rangle$

lemma *Cauchy-converges-subseq*:
fixes $u:\text{nat} \Rightarrow 'a$
assumes *Cauchy u*
strict-mono r
 $(u \circ r) \longrightarrow l$
shows $u \longrightarrow l$
 $\langle proof \rangle$

end

3.1 The Construction with Classes

class *restriction-σ* = *restriction-space* +
fixes *restriction-σ* :: $\langle 'a \text{ itself} \Rightarrow \text{nat} \Rightarrow \text{real} \rangle (\sigma_\downarrow)$

$\langle ML \rangle$

```

class non-decseq-restriction-space =
  uniformity-dist + open-uniformity + restriction- $\sigma$  +
  — We do not assume the restriction sequence to be decseq yet.
  assumes restriction- $\sigma$ -tendsto-zero' :  $\langle \sigma_{\downarrow} \text{TYPE('a)} \longrightarrow 0 \rangle$ 
  and zero-less-restriction- $\sigma'$  [simp] :  $\langle 0 < \sigma_{\downarrow} \text{TYPE('a)} n \rangle$ 

  and dist-restriction-is' :  $\langle \text{dist } x y = (\text{INF } n \in \{n. x \downarrow n = y \downarrow n\}.$ 
   $\sigma_{\downarrow} \text{TYPE('a)} n) \rangle$ 
begin

  sublocale NonDecseqRestrictionSpace  $\langle (\downarrow) \rangle$   $\langle (=) \rangle$   $\langle \text{UNIV} \rangle$   $\langle \sigma_{\downarrow} \text{TYPE}('a) \rangle$  dist
   $\langle \text{proof} \rangle$ 

end

⟨ML⟩

We hide duplicated facts  $\sigma_{\downarrow} \text{TYPE}(?'a) \longrightarrow 0$ 
 $0 < \sigma_{\downarrow} \text{TYPE}(?'a) ?n$ 
 $\text{dist } ?x ?y = \text{Inf } (\text{restriction-}\sigma\text{-related-set } ?x ?y).$ 
hide-fact restriction- $\sigma$ -tendsto-zero' zero-less-restriction- $\sigma'$  dist-restriction-is'

context non-decseq-restriction-space begin

  subclass ultrametric-space
   $\langle \text{proof} \rangle$ 

end

context non-decseq-restriction-space begin

  lemma restriction-tendsto-self :  $\langle (\lambda n. x \downarrow n) \longrightarrow x \rangle$ 
   $\langle \text{proof} \rangle$ 

end

class decseq-restriction-space = non-decseq-restriction-space +
  assumes decseq-restriction- $\sigma'$  :  $\langle \text{decseq } (\sigma_{\downarrow} \text{TYPE('a)}) \rangle$ 
begin

  sublocale DecseqRestrictionSpace  $\langle (\downarrow) \rangle$   $\langle (=) \rangle$   $\langle \text{UNIV} :: 'a \text{ set} \rangle$ 
   $\langle \text{restriction-}\sigma \text{ TYPE('a)} \rangle$  dist

```

```

⟨proof⟩
lemmas dist-restriction-is-bis-simplified = dist-restriction-is-bis[simplified]
and not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified
=
not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq[simplified]

end

```

We hide duplicated fact *decseq* ($\sigma \downarrow \text{TYPE}(\text{'a})$).

```
hide-fact decseq-restriction-σ'
```

```

class strict-decseq-restriction-space = non-decseq-restriction-space +
assumes strict-decseq-restriction-σ : ⟨strict-decseq ( $\sigma \downarrow \text{TYPE}(\text{'a})$ )⟩
begin

subclass decseq-restriction-space
⟨proof⟩

end

```

Generic Properties

```
lemma (in metric-space) dist-sequences-tendsto-zero-imp-tendsto-iff :
⟨(λn. dist ( $\sigma n$ ) ( $\psi n$ )) —→ 0 ⟹ σ —→ Σ ↔ ψ —→ Σ⟩
⟨proof⟩
```

```
lemma (in non-decseq-restriction-space) restricted-sequence-tendsto-iff
:
⟨(λn. σ n ↓ n) —→ Σ ↔ σ —→ Σ⟩
⟨proof⟩
```

```
lemma (in non-decseq-restriction-space) Cauchy-restriction-chain :
⟨Cauchy σ⟩ if ⟨chain↓ σ⟩
⟨proof⟩
```

```
lemma (in non-decseq-restriction-space) restriction-tendsto-imp-tendsto
:
⟨σ —→ Σ⟩ if ⟨σ —→ Σ⟩
⟨proof⟩
```

In Decseq Restriction Space

```
context decseq-restriction-space begin
```

```

lemma le-dist-to-restriction-eqE :
  obtains k where ⟨n ≤ k⟩ ⟨ $\bigwedge x y :: 'a.$  dist x y ≤  $\sigma \downarrow$  TYPE('a) k
   $\implies x \downarrow n = y \downarrow n$ ⟩
  ⟨proof⟩

```

```

theorem tendsto-iff-restriction-tendsto : ⟨ $\sigma \longrightarrow \Sigma \longleftrightarrow \sigma \dashrightarrow \Sigma$ ⟩
  ⟨proof⟩

```

```

corollary convergent-iff-restriction-convergent : ⟨convergent  $\sigma \longleftrightarrow$ 
  convergent $\downarrow \sigma$ ⟩
  ⟨proof⟩

```

```

theorem complete-iff-restriction-complete :
  ⟨( $\forall \sigma.$  Cauchy  $\sigma \longrightarrow$  convergent  $\sigma$ )  $\longleftrightarrow$  ( $\forall \sigma.$  chain $\downarrow \sigma \longrightarrow$  convergent $\downarrow \sigma$ )⟩
  — The following result shows that we have not lost anything with
  our definitions of convergence, completeness, etc. specific to restriction
  spaces.
  ⟨proof⟩

```

end

The following classes will be useful later.

```

class complete-decseq-restriction-space = decseq-restriction-space +
  assumes restriction-chain-imp-restriction-convergent' : ⟨chain $\downarrow \sigma \implies$  convergent $\downarrow \sigma$ ⟩
  begin

```

```

    subclass complete-restriction-space
    ⟨proof⟩

```

```

    subclass complete-ultrametric-space
    ⟨proof⟩

```

end

We hide duplicated fact $chain \downarrow ?\sigma \implies convergent \downarrow ?\sigma$.

hide-fact restriction-chain-imp-restriction-convergent'

```

class complete-strict-decseq-restriction-space = strict-decseq-restriction-space
+
  assumes restriction-chain-imp-restriction-convergent' : ⟨chain $\downarrow \sigma \implies$  convergent $\downarrow \sigma$ ⟩

```

```

begin

subclass complete-decseq-restriction-space
  ⟨proof⟩

end

We hide duplicated fact  $chain_{\downarrow} ?\sigma \implies convergent_{\downarrow} ?\sigma$ .
hide-fact restriction-chain-imp-restriction-convergent'

class restriction- $\delta$  = restriction- $\sigma$  +
  fixes restriction- $\delta$  :: ⟨'a itself  $\Rightarrow$  real⟩ ⟨ $\delta_{\downarrow}$ ⟩
  assumes zero-less-restriction- $\delta$  [simp] : ⟨ $0 < \delta_{\downarrow} \text{TYPE('a)}$ ⟩
    and restriction- $\delta$ -less-one [simp] : ⟨ $\delta_{\downarrow} \text{TYPE('a)} < 1$ ⟩
begin

lemma zero-le-restriction- $\delta$  [simp] : ⟨ $0 \leq \delta_{\downarrow} \text{TYPE('a)}$ ⟩
  and restriction- $\delta$ -le-one [simp] : ⟨ $\delta_{\downarrow} \text{TYPE('a)} \leq 1$ ⟩
  and zero-le-pow-restriction- $\delta$  [simp] : ⟨ $0 \leq \delta_{\downarrow} \text{TYPE('a)} \wedge n$ ⟩
  ⟨proof⟩

lemma pow-restriction- $\delta$ -le-one [simp] : ⟨ $\delta_{\downarrow} \text{TYPE('a)} \wedge n \leq 1$ ⟩
  ⟨proof⟩

lemma pow-restriction- $\delta$ -less-one [simp] : ⟨ $n \neq 0 \implies \delta_{\downarrow} \text{TYPE('a)} \wedge n < 1$ ⟩
  ⟨proof⟩

end

⟨ML⟩

class at-least-geometric-restriction-space =
  uniformity-dist + open-uniformity + restriction- $\delta$  +
  assumes zero-less-restriction- $\sigma'$  : ⟨ $0 < \sigma_{\downarrow} \text{TYPE('a)} n$ ⟩
    and restriction- $\sigma$ -le :
      ⟨ $\sigma_{\downarrow} \text{TYPE('a)} (\text{Suc } n) \leq \delta_{\downarrow} \text{TYPE('a)} * \sigma_{\downarrow} \text{TYPE('a)} n$ ⟩
    and dist-restriction-is' :
      ⟨dist x y = (INF n ∈ {n. x ↓ n = y ↓ n}. σ_{\downarrow} \text{TYPE('a)} n)⟩

⟨ML⟩

context at-least-geometric-restriction-space begin

lemma restriction- $\sigma$ -le-restriction- $\sigma$ -times-pow-restriction- $\delta$  :
  ⟨ $\sigma_{\downarrow} \text{TYPE('a)} (n + k) \leq \sigma_{\downarrow} \text{TYPE('a)} n * \delta_{\downarrow} \text{TYPE('a)} \wedge k$ ⟩

```

$\langle proof \rangle$

```
lemma restriction- $\sigma$ -le-pow-restriction- $\delta$  :
   $\langle \sigma_\downarrow \text{TYPE}('a) n \leq \sigma_\downarrow \text{TYPE}('a) 0 * \delta_\downarrow \text{TYPE}('a) \wedge n \rangle$ 
   $\langle proof \rangle$ 
```

```
subclass strict-decseq-restriction-space
   $\langle proof \rangle$ 
```

```
lemma  $\langle 0 < \delta_\downarrow \text{TYPE}('a) \wedge n \rangle \langle proof \rangle$ 
```

end

We hide duplicated facts $0 < \sigma_\downarrow \text{TYPE}('a) ?n$
 $dist ?x ?y = Inf (\text{restriction-}\sigma\text{-related-set} ?x ?y)$.

hide-fact zero-less-restriction- σ' dist-restriction-is'

```
class complete-at-least-geometric-restriction-space = at-least-geometric-restriction-space
```

```
+  
  assumes restriction-chain-imp-restriction-convergent' :  $\langle \text{chain}_\downarrow \sigma$   

 $\implies \text{convergent}_\downarrow \sigma \rangle$   

begin
```

```
subclass complete-strict-decseq-restriction-space
   $\langle proof \rangle$ 
```

end

We hide duplicated fact $\text{chain}_\downarrow ?\sigma \implies \text{convergent}_\downarrow ?\sigma$.

hide-fact restriction-chain-imp-restriction-convergent'

$\langle ML \rangle$

```
class geometric-restriction-space = uniformity-dist + open-uniformity  

+ restriction- $\delta$  +
```

```
  assumes restriction- $\sigma$ -is :  $\langle \sigma_\downarrow \text{TYPE}('a) n = \delta_\downarrow \text{TYPE}('a) \wedge n \rangle$   

    and dist-restriction-is' :  $\langle \text{dist } x y = (\text{INF } n \in \{n. x_\downarrow n = y_\downarrow n\}.$   

 $\sigma_\downarrow \text{TYPE}('a) n) \rangle$ 
```

begin

This is what “restriction space” usually mean in the literature. The previous classes are generalizations of this concept (even this one is a generalization, since we usually have $\delta \downarrow \text{TYPE}('a) = 1 / 2$).

subclass *at-least-geometric-restriction-space*
⟨proof⟩

lemma *⟨0 < δ↓ TYPE('a) ^ n⟩* *⟨proof⟩*

end

⟨ML⟩

We hide duplicated fact *dist ?x ?y = Inf (restriction-σ-related-set ?x ?y)*.

hide-fact *dist-restriction-is'*

class *complete-geometric-restriction-space = geometric-restriction-space*
+
assumes *restriction-chain-imp-restriction-convergent' : ⟨chain↓ σ*
 $\implies \text{convergent}_\downarrow \sigma$
begin

subclass *complete-at-least-geometric-restriction-space*
⟨proof⟩

end

We hide duplicated fact *chain↓ σ ⇒ convergent↓ σ*.

hide-fact *restriction-chain-imp-restriction-convergent'*

theorem *geometric-restriction-space-completeI : ⟨convergent σ*
if *⟨ʃ σ :: nat ⇒ 'a. restriction-chain σ ⇒ ∃ Σ. ∀ n. Σ ↓ n = σ n⟩*
and *⟨Cauchy σ⟩* **for** *σ :: ⟨nat ⇒ 'a :: geometric-restriction-space⟩*
⟨proof⟩
lemma **(in** *non-decseq-restriction-space*) *restriction-cball-subset-cball*
 $:$
⟨σ↓ TYPE('a) n ≤ r ⇒ B↓(Σ, n) ⊆ {x. dist Σ x ≤ r}⟩
⟨proof⟩

corollary *restriction-cball-subset-cball-bis :*
⟨σ↓ TYPE('a) n ≤ r ⇒ B↓(Σ, n) ⊆ cball Σ r⟩
for *Σ :: ⟨'a :: non-decseq-restriction-space⟩*

$\langle proof \rangle$

lemma (in non-decseq-restriction-space) restriction-ball-subset-ball :
 $\langle \sigma_{\downarrow} \text{TYPE('a)} n < r \implies \text{restriction-ball } \Sigma n \subseteq \{x. \text{dist } \Sigma x < r\} \rangle$
 $\langle proof \rangle$

corollary restriction-ball-subset-ball-bis :
 $\langle \sigma_{\downarrow} \text{TYPE('a)} n < r \implies \text{restriction-ball } \Sigma n \subseteq \text{ball } \Sigma r \rangle$
for $\Sigma :: \langle 'a :: \text{non-decseq-restriction-space} \rangle$
 $\langle proof \rangle$

lemma (in strict-decseq-restriction-space)
restriction-cball-is-cball : $\langle \mathcal{B}_{\downarrow}(\Sigma, n) = \{x. \text{dist } \Sigma x \leq \sigma_{\downarrow} \text{TYPE('a)} n\} \rangle$
 $\langle proof \rangle$

lemma restriction-cball-is-cball-bis : $\langle \mathcal{B}_{\downarrow}(\Sigma, n) = \text{cball } \Sigma (\sigma_{\downarrow} \text{TYPE('a)} n) \rangle$
for $\Sigma :: \langle 'a :: \text{strict-decseq-restriction-space} \rangle$
 $\langle proof \rangle$

lemma (in strict-decseq-restriction-space)
restriction-ball-is-ball : $\langle \text{restriction-ball } \Sigma n = \{x. \text{dist } \Sigma x < \sigma_{\downarrow} \text{TYPE('a)} n\} \rangle$
 $\langle proof \rangle$

lemma restriction-ball-is-ball-bis : $\langle \text{restriction-ball } \Sigma n = \text{ball } \Sigma (\sigma_{\downarrow} \text{TYPE('a)} n) \rangle$
for $\Sigma :: \langle 'a :: \text{strict-decseq-restriction-space} \rangle$
 $\langle proof \rangle$

lemma isCont-iff-restriction-cont-at : $\langle \text{isCont } f \Sigma \longleftrightarrow \text{restriction-cont-at } f \Sigma \rangle$
for $f :: \langle 'a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{decseq-restriction-space} \rangle$
 $\langle proof \rangle$

lemma (in strict-decseq-restriction-space)
open-iff-restriction-open : $\langle \text{open } U \longleftrightarrow \text{open}_{\downarrow} U \rangle$
 $\langle proof \rangle$

```

lemma (in strict-decseq-restriction-space)
  closed-iff-restriction-closed : <closed U  $\longleftrightarrow$  closed↓ U>
  ⟨proof⟩

```

```

lemma continuous-on-iff-restriction-cont-on :
  ⟨open U  $\implies$  continuous-on U f  $\longleftrightarrow$  restriction-cont-on f U>
  for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: decseq-restriction-space⟩
  ⟨proof⟩

```

3.2 Equivalence between Lipschitz Map and Restriction shift Map

For a function $f: 'a \Rightarrow 'b$, it is equivalent to have *restriction-shift-on* $f k A$ and *lipschitz-with-on* $f (\delta_{\downarrow} \text{TYPE}('b))^k A$ when ' a ' is of sort *geometric-restriction-space* and $\sigma_{\downarrow} \text{TYPE}('b) = \sigma_{\downarrow} \text{TYPE}('a)$ (' b ' is then necessarily also of sort *geometric-restriction-space*).

Weaker versions of this result can be established with weaker assumptions on the sort, this is what we do below.

```

lemma restriction-shift-nonneg-imp-lipschitz-with-on :
  ⟨lipschitz-with-on f (restriction-δ TYPE('b)  $\wedge$  k) A⟩ if ⟨restriction-shift-on f (int k) A⟩
  and le-restriction-σ : ⟨ $\bigwedge n. \text{restriction-}\sigma \text{TYPE}('b) n \leq \text{restriction-}\sigma \text{TYPE}('a) n$ ⟩
  for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
  ⟨proof⟩

```

```

corollary restriction-shift-nonneg-imp-lipschitz-with :
  ⟨[restriction-shift f (int k);  $\bigwedge n. \text{restriction-}\sigma \text{TYPE}('b) n \leq \text{restriction-}\sigma \text{TYPE}('a) n]$ ]
   $\implies$  lipschitz-with f (restriction-δ TYPE('b)  $\wedge$  k)⟩
  for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
  ⟨proof⟩

```

```

lemma lipschitz-with-on-imp-restriction-shift-neg-on :
  ⟨restriction-shift-on f (- int k) A⟩ if ⟨lipschitz-with-on f (restriction-δ TYPE('b) powi - int k) A⟩
  and le-restriction-σ : ⟨ $\bigwedge n. \text{restriction-}\sigma \text{TYPE}('a) n \leq \text{restriction-}\sigma \text{TYPE}('b) n$ ⟩
  for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
  ⟨proof⟩

```

```

corollary lipschitz-with-imp-restriction-shift-neg :

```

```

⟨[lipschitz-with f (restriction- $\delta$  TYPE('b) powi – int k);
  ∧n. restriction- $\sigma$  TYPE('a) n ≤ restriction- $\sigma$  TYPE('b) n]
  ⇒ restriction-shift f (– int k))⟩
for f :: ⟨'a :: decseq-restriction-space ⇒ 'b :: at-least-geometric-restriction-space⟩
⟨proof⟩

```

We obtained that *restriction-shift* implies *lipschitz-with* when $0 \leq k$ and that *lipschitz-with* implies *restriction-shift* when $k \leq 0$.

To cover the remaining cases, we give move from *at-least-geometric-restriction-space* to *geometric-restriction-space*.

theorem *lipschitz-with-on-iff-restriction-shift-on* :

```

⟨lipschitz-with-on f (restriction- $\delta$  TYPE('b) powi k) A ↔ restriction-shift-on f k A⟩
if same-restriction- $\sigma$  : ⟨restriction- $\sigma$  TYPE('b) = restriction- $\sigma$  TYPE('a)⟩
for f :: ⟨'a :: decseq-restriction-space ⇒ 'b :: geometric-restriction-space⟩
⟨proof⟩

```

corollary *lipschitz-with-iff-restriction-shift* :

```

⟨restriction- $\sigma$  TYPE('b) = restriction- $\sigma$  TYPE('a) ⇒
  lipschitz-with f (restriction- $\delta$  TYPE('b) powi k) ↔ restriction-shift
  f k⟩
for f :: ⟨'a :: decseq-restriction-space ⇒ 'b :: geometric-restriction-space⟩
⟨proof⟩

```

4 Functions

4.1 Restriction Space

instantiation ⟨fun⟩ :: (type, restriction- σ) restriction- σ
begin

definition restriction- σ -fun :: ⟨('a ⇒ 'b) itself ⇒ nat ⇒ real⟩
where ⟨restriction- σ -fun - ≡ restriction- σ TYPE('b)⟩

instance ⟨proof⟩

end

instantiation ⟨fun⟩ :: (type, non-decseq-restriction-space) non-decseq-restriction-space
begin

definition dist-fun :: ⟨['a ⇒ 'b, 'a ⇒ 'b] ⇒ real⟩
where ⟨dist-fun fg ≡ INF n ∈ restriction-related-set fg. restriction- σ TYPE('a ⇒ 'b) n⟩

definition uniformity-fun :: ⟨((‘a ⇒ ‘b) × (‘a ⇒ ‘b)) filter⟩
where ⟨uniformity-fun ≡ INF e ∈ {0 <..}. principal {(x, y). dist x y < e}⟩

```

definition open-fun :: <('a ⇒ 'b) set ⇒ bool>
  where <open-fun U ≡ ∀x∈U. eventually (λ(x', y). x' = x → y ∈ U) uniformity>

instance ⟨proof⟩

end

instance ⟨fun⟩ :: (type, decseq-restriction-space) decseq-restriction-space
⟨proof⟩

instance ⟨fun⟩ :: (type, strict-decseq-restriction-space) strict-decseq-restriction-space
⟨proof⟩

instantiation ⟨fun⟩ :: (type, restriction-δ) restriction-δ
begin

definition restriction-δ-fun :: <('a ⇒ 'b) itself ⇒ real>
  where <restriction-δ-fun - ≡ restriction-δ TYPE('b)>

instance ⟨proof⟩

end

instance ⟨fun⟩ :: (type, at-least-geometric-restriction-space) at-least-geometric-restriction-space
⟨proof⟩

instance ⟨fun⟩ :: (type, geometric-restriction-space) geometric-restriction-space
⟨proof⟩

lemma dist-image-le-dist-fun : <dist (f x) (g x) ≤ dist f g>
  for f g :: <'a ⇒ 'b :: non-decseq-restriction-space>
⟨proof⟩

lemma Sup-dist-image-le-dist-fun : <(SUP x. dist (f x) (g x)) ≤ dist f g>
  for f g :: <'a ⇒ 'b :: non-decseq-restriction-space>
⟨proof⟩

context fixes f g :: <'a ⇒ 'b :: decseq-restriction-space> begin

lemma reached-dist-fun : <∃x. dist f g = dist (f x) (g x)>

```

$\langle proof \rangle$

lemma *dist-fun-eq-Sup-dist-image* : $\langle dist f g = (SUP x. dist (f x) (g x)) \rangle$
 $\langle proof \rangle$

lemma *fun-restriction-space-Sup-properties* :
 $\langle dist (f x) (g x) \leq dist f g \rangle$
 $\langle (\forall x. dist (f x) (g x) \leq b) \implies dist f g \leq b \rangle$
 $\langle proof \rangle$

end

4.2 Completeness

Actually we can obtain even better: when the instance '*b*' of *decseq-restriction-space* is also an instance of *complete-space*, the type '*a* \Rightarrow '*b*' is an instance of *complete-space*.

This is because when '*b*' is an instance of *decseq-restriction-space* (and not only *non-decseq-restriction-space*) the distance between two functions is reached (see $\exists x. dist ?f ?g = dist (?f x) (?g x)$).

The only remaining thing is to prove that completeness is preserved on higher-order.

instance *fun* :: (*type, complete-decseq-restriction-space*) *complete-decseq-restriction-space*
 $\langle proof \rangle$

instance *fun* :: (*type, complete-strict-decseq-restriction-space*) *complete-strict-decseq-restriction-space*
 $\langle proof \rangle$

instance *fun* :: (*type, complete-at-least-geometric-restriction-space*) *complete-at-least-geometric-restriction-space*
 $\langle proof \rangle$

instance *fun* :: (*type, complete-geometric-restriction-space*) *complete-geometric-restriction-space*
 $\langle proof \rangle$

4.3 Kind of Extensionality

context *fixes f* :: $\langle [a :: metric-space, b :: type] \Rightarrow c :: decseq-restriction-space \rangle$ **begin**

lemma *lipschitz-with-simplification*:
 $\langle lipschitz-with f \alpha \longleftrightarrow (\forall y. lipschitz-with (\lambda x. f x y) \alpha) \rangle$

```
 $\langle proof \rangle$ 
```

```
lemma non-expanding-simplification :  
  ‹non-expanding f  $\longleftrightarrow$  ( $\forall y$ . non-expanding ( $\lambda x$ . f x y))›  
 $\langle proof \rangle$ 
```

```
lemma contraction-with-simplification:  
  ‹contraction-with f  $\alpha$   $\longleftrightarrow$  ( $\forall y$ . contraction-with ( $\lambda x$ . f x y)  $\alpha$ )›  
 $\langle proof \rangle$ 
```

```
end
```

5 Product

The product type ' $a \times b$ ' of two metric spaces is already instantiated as a metric space by setting $dist x y = sqrt((dist(fst x)(fst y))^2 + (dist(snd x)(snd y))^2)$. Unfortunately, this definition is not compatible with the distance required by the *non-decseq-restriction-space*. We first have to define a new product type with a trivial **typedef**.

5.1 Isomorphic Product Construction

5.1.1 Definition and First Properties

```
typedef ('a, 'b) prodmax ((- ×max / -) [21, 20] 20) = ‹UNIV :: ('a  
  × 'b) set›  
morphisms from-prodmax to-prodmax  $\langle proof \rangle$ 
```

```
declare from-prodmax-inject [simp]  
from-prodmax-inverse [simp]
```

```
lemmas to-prodmax-inject-simplified [simp] = to-prodmax-inject [simplified]  
and to-prodmax-inverse-simplified [simp] = to-prodmax-inverse [simplified]
```

```
lemmas to-prodmax-induct-simplified = to-prodmax-induct [simplified]  
and to-prodmax-cases-simplified = to-prodmax-cases [simplified]  
and from-prodmax-induct-simplified = from-prodmax-induct [simplified]  
and from-prodmax-cases-simplified = from-prodmax-cases [simplified]
```

```
setup-lifting type-definition-prodmax
```

lift-definition $\text{Pair}_{\max} :: \langle 'a \Rightarrow 'b \Rightarrow 'a \times_{\max} 'b \rangle$ **is** $\text{Pair} \langle \text{proof} \rangle$

free-constructors case-prod_{\max} **for** $\text{Pair}_{\max} \text{fst}_{\max} \text{snd}_{\max}$
 $\langle \text{proof} \rangle$

lemma $\text{fst}_{\max}\text{-def} : \langle \text{fst}_{\max} \equiv \text{map-fun from-prod}_{\max} \text{id fst} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{fst}_{\max}\text{-rep-eq} : \langle \text{fst}_{\max} x = \text{fst} (\text{from-prod}_{\max} x) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{fst}_{\max}\text{-abs-eq [simp]} : \langle \text{fst}_{\max} (\text{to-prod}_{\max} y) = \text{fst} y \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{fst}_{\max}\text{-transfer [transfer-rule]} : \langle \text{rel-fun (pcr-prod}_{\max} (=) (=))$
 $(=) \text{fst fst}_{\max} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{snd}_{\max}\text{-def} : \langle \text{snd}_{\max} \equiv \text{map-fun from-prod}_{\max} \text{id snd} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{snd}_{\max}\text{-rep-eq} : \langle \text{snd}_{\max} x = \text{snd} (\text{from-prod}_{\max} x) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{snd}_{\max}\text{-abs-eq [simp]} : \langle \text{snd}_{\max} (\text{to-prod}_{\max} y) = \text{snd} y \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{snd}_{\max}\text{-transfer [transfer-rule]} : \langle \text{rel-fun (pcr-prod}_{\max} (=)$
 $(=) (=) \text{snd snd}_{\max} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{case-prod}_{\max}\text{-def} : \langle \text{case-prod}_{\max} \equiv \text{map-fun id (map-fun}$
 $\text{from-prod}_{\max} \text{id}) \text{ case-prod} \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{case-prod}_{\max}\text{-rep-eq} : \langle \text{case-prod}_{\max} f p = (\text{case from-prod}_{\max}$
 $p \text{ of } (x, y) \Rightarrow f x y) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{case-prod}_{\max}\text{-abs-eq [simp]} : \langle \text{case-prod}_{\max} f (\text{to-prod}_{\max} q)$
 $= (\text{case q of } (x, y) \Rightarrow f x y) \rangle$
 $\langle \text{proof} \rangle$

lemma $\text{case-prod}_{\max}\text{-transfer [transfer-rule]} : \langle \text{rel-fun (=) (rel-fun}$
 $(\text{pcr-prod}_{\max} (=) (=) (=)) \text{ case-prod case-prod}_{\max} \rangle$
 $\langle \text{proof} \rangle$

5.2 Syntactic Sugar

The following syntactic sugar is of course recovered from the theory *HOL.Product-Type*.

```
nonterminal tuple-argsmax and patternsmax
syntax
-tuplemax :: 'a ⇒ tuple-argsmax ⇒ 'a ×max 'b      ((1⟨-, / -⟩)⟨)
-tuple-argmax :: 'a ⇒ tuple-argsmax                  (↔)
-tuple-argsmax :: 'a ⇒ tuple-argsmax ⇒ tuple-argsmax  (⟨-, / -⟩)
-patternmax   :: pttrn ⇒ patternsmax ⇒ pttrn       ((⟨-, / -⟩)⟨)
              :: pttrn ⇒ patternsmax                   (↔)
-patternsmax :: pttrn ⇒ patternsmax ⇒ patternsmax  (⟨-, / -⟩)
translations
⟨x, y⟩ ⇌ CONST Pairmax x y
-patternmax x y ⇌ CONST Pairmax x y
-patternsmax x y ⇌ CONST Pairmax x y
-tuplemax x (-tuple-argsmax y z) ⇌ -tuplemax x (-tuple-argmax
(-tuplemax y z))
λ⟨x, y, zs⟩. b ⇌ CONST case-prodmax (λx ⟨y, zs⟩. b)
λ⟨x, y⟩. b ⇌ CONST case-prodmax (λx y. b)
-abs (CONST Pairmax x y) t → λ⟨x, y⟩. t
— This rule accommodates tuples in case C ... ⟨x, y⟩ ... ⇒ ...:
The ⟨x, y⟩ is parsed as Pairmax x y because it is logic, not pttrn.
```

With this syntactic sugar, one can write *case a of* ⟨b, c, d, e⟩ ⇒ ⟨c, d⟩, λ⟨y, u⟩. a, λ⟨a, b⟩. ⟨a, b, c, d, e⟩, λ⟨a, b, c⟩. a, ... as for the type 'a × 'b'.

```
lemmas to-prodmax-tuple [simp] = Pairmax.abs-eq[symmetric]
and from-prodmax-tuplemax [simp] = Pairmax.rep-eq
```

5.3 Product

We first redo the work of *Restriction-Spaces.Restriction-Spaces-Prod*.

```
instantiation prodmax :: (restriction, restriction) restriction
begin
```

```
lift-definition restriction-prodmax :: ⟨'a ×max 'b ⇒ nat ⇒ 'a ×max
'b⟩ is ⟨(↓)⟩ ⟨proof⟩
```

```
lemma restriction-prodmax-def' : ⟨p ↓ n = ⟨fstmax p ↓ n, sndmax p
↓ n⟩⟩
⟨proof⟩
```

```
instance ⟨proof⟩
```

```
end
```

```

instance prodmax :: (restriction-space, restriction-space) restriction-space
  ⟨proof⟩

instantiation prodmax :: (restriction-σ, restriction-σ) restriction-σ
begin

definition restriction-σ-prodmax :: ⟨('a ×max 'b) itself ⇒ nat ⇒ real⟩
  where ⟨restriction-σ-prodmax - n ≡
    max (restriction-σ TYPE('a) n) (restriction-σ TYPE('b) n)⟩

instance ⟨proof⟩
end

instantiation prodmax :: (non-decseq-restriction-space, non-decseq-restriction-space)
  non-decseq-restriction-space
begin

definition dist-prodmax :: ⟨['a ×max 'b, 'a ×max 'b] ⇒ real⟩
  where ⟨dist-prodmax f g ≡ INF n ∈ restriction-related-set f g. re-
    striction-σ TYPE('a ×max 'b) n⟩

definition uniformity-prodmax :: ⟨((('a ×max 'b) × 'a ×max 'b) filter)
  where ⟨uniformity-prodmax ≡ INF e ∈ {0 <..}. principal {(x, y). dist
    x y < e}⟩

definition open-prodmax :: ⟨('a ×max 'b) set ⇒ bool⟩
  where ⟨open-prodmax U ≡ ∀ x ∈ U. eventually (λ(x', y). x' = x →
    y ∈ U) uniformity⟩

instance
  ⟨proof⟩

end

instance prodmax :: (decseq-restriction-space, decseq-restriction-space)
  decseq-restriction-space
  ⟨proof⟩

instance prodmax :: (strict-decseq-restriction-space, strict-decseq-restriction-space)
  strict-decseq-restriction-space
  ⟨proof⟩

instantiation prodmax :: (restriction-δ, restriction-δ) restriction-δ
begin

```

```

definition restriction- $\delta$ -prodmax :: <('a ×max 'b) itself ⇒ real>
  where <restriction- $\delta$ -prodmax - ≡ max (restriction- $\delta$  TYPE('a))
    (restriction- $\delta$  TYPE('b))>

instance ⟨proof⟩

end

instance prodmax :: (at-least-geometric-restriction-space, at-least-geometric-restriction-space)
  at-least-geometric-restriction-space
  ⟨proof⟩

instance prodmax :: (geometric-restriction-space, geometric-restriction-space)
  geometric-restriction-space
  ⟨proof⟩

```

```

lemma max-dist-projections-le-dist-prodmax :
  <max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1) (sndmax p2)))
  ≤ dist p1 p2>
  ⟨proof⟩

```

5.4 Completeness

5.4.1 Preliminaries

default-sort non-decseq-restriction-space — Otherwise we should always specify.

```

lemma restriction- $\sigma$ -prodmax-commute :
  <restriction- $\sigma$  TYPE('b ×max 'a) = restriction- $\sigma$  TYPE('a ×max 'b)>
  ⟨proof⟩

```

```

definition dist-left-prodmax :: <('a ×max 'b) itself ⇒ 'a ⇒ 'a ⇒ real>
  where <dist-left-prodmax - x y ≡ INF n ∈ restriction-related-set x
    y. restriction- $\sigma$  TYPE('a ×max 'b) n>

```

```

definition dist-right-prodmax :: <('a ×max 'b) itself ⇒ 'b ⇒ 'b ⇒ real>
  where <dist-right-prodmax - x y ≡ INF n ∈ restriction-related-set x
    y. restriction- $\sigma$  TYPE('a ×max 'b) n>

```

```

lemma dist-right-prodmax-is-dist-left-prodmax :
  <dist-right-prodmax TYPE('b ×max 'a) = dist-left-prodmax TYPE('a
  ×max 'b)>
  ⟨proof⟩

```

```
lemma dist-le-dist-left-prodmax : <dist x y ≤ dist-left-prodmax TYPE('a ×max 'b) x y>
⟨proof⟩
```

```
lemma dist-le-dist-right-prodmax : <dist x y ≤ dist-right-prodmax TYPE('b ×max 'a) x y>
⟨proof⟩
```

```
lemma
  fixes p1 p2 :: <'a :: decseq-restriction-space ×max 'b :: decseq-restriction-space>
  shows dist-prodmax-le-max-dist-left-prodmax-dist-right-prodmax :
    <dist p1 p2 ≤ max (dist-left-prodmax TYPE('a ×max 'b) (fstmax p1) (fstmax p2)) +
      (dist-right-prodmax TYPE('a ×max 'b) (sndmax p1) (sndmax p2))>
  ⟨proof⟩
```

default-sort type — Back to normal.

5.4.2 Complete Restriction Space

When the instances '*a*' and '*b*' of *decseq-restriction-space* are also instances of *complete-space*, the type '*a* ×_{max} '*b*' is an instance of *complete-space*.

```
lemma restriction-chain-prodmax-iff :
  <restriction-chain σ ←→ restriction-chain (λn. fstmax (σ n)) ∧
    restriction-chain (λn. sndmax (σ n))>
  ⟨proof⟩
```

```
lemma restriction-tendsto-prodmax-iff :
  <σ → Σ ←→ (λn. fstmax (σ n)) → fstmax Σ ∧ (λn. sndmax (σ n)) → sndmax Σ>
  ⟨proof⟩
```

```
lemma restriction-convergent-prodmax-iff :
  <restriction-convergent σ ←→ restriction-convergent (λn. fstmax (σ n)) ∧
    restriction-convergent (λn. sndmax (σ n))>
  ⟨proof⟩
```

```
instance prodmax :: (complete-decseq-restriction-space, complete-decseq-restriction-space)
  complete-decseq-restriction-space
  ⟨proof⟩
```

```

instance prodmax :: (complete-strict-decseq-restriction-space, complete-strict-decseq-restriction-space)
  complete-strict-decseq-restriction-space
  ⟨proof⟩

instance prodmax :: (complete-at-least-geometric-restriction-space, complete-at-least-geometric-restriction-space)
  complete-at-least-geometric-restriction-space
  ⟨proof⟩

instance prodmax :: (complete-geometric-restriction-space, complete-geometric-restriction-space)
  complete-geometric-restriction-space
  ⟨proof⟩

```

When the types ' a ' and ' b ' share the same restriction sequence, we have the following equality.

```

lemma same-restriction-σ-imp-restriction-σ-prodmax-is [simp] :
  <restriction-σ TYPE('b :: non-decseq-restriction-space) =
    restriction-σ TYPE('a :: non-decseq-restriction-space) ⟹
    restriction-σ TYPE('a ×max 'b) = restriction-σ TYPE('a),
  ⟨proof⟩

```

```

lemma same-restriction-σ-imp-dist-prodmax-eq-max-dist-projections :
  <dist p1 p2 = max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1)
  (sndmax p2)))
  if same-restriction-σ [simp] : <restriction-σ TYPE('b) = restriction-σ
  TYPE('a)
  for p1 p2 :: <'a :: decseq-restriction-space ×max 'b :: decseq-restriction-space>
  ⟨proof⟩

```

Now, one can write things like $v \langle x, y \rangle . f \langle x, y \rangle$.

We could of course imagine more support for ' $a ×_{max} b$ ' type, but as restriction spaces are intended to be used without recourse to metric spaces, we have not undertaken this task for the time being.

6 Main entry Point