

# Providing restriction Spaces with an ultrametric Structure

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## Abstract

We investigate the relationship between restriction spaces and classical metric structures by instantiating the former as ultrametric spaces. This is classically captured by defining the distance as

$$\text{dist } x \ y = \inf_{x \downarrow n = y \downarrow n} \left( \frac{1}{2} \right)^n$$

but we actually generalize this perspective by introducing a hierarchy of increasingly refined type classes to systematically relate ultrametric and restriction-based notions. This layered approach enables a precise comparison of structural and topological properties. In the end, our main result establishes that completeness in the sense of restriction spaces coincides with standard metric completeness, thus bridging the gap between `Restriction_Spaces` and Banach's fixed-point theorem established in `HOL-Analysis`.

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## 1 Definitions on Functions of Metric Space

In this theory, we define the notion of lipschitz map, non-expanding map and contraction map. We also establish correspondences.

### 1.1 Definitions

#### 1.1.1 Lipschitz Map

This notion is a generalization of contraction map and non-expanding map.

**definition** *lipschitz-with-on* ::  $\langle [a :: metric-space \Rightarrow b :: metric-space, real, 'a set] \Rightarrow bool \rangle$   
**where**  $\langle lipschitz-with-on f \alpha A \equiv 0 \leq \alpha \wedge (\forall x \in A. \forall y \in A. dist(f x) (f y) \leq \alpha * dist x y) \rangle$

**abbreviation** *lipschitz-with* ::  $\langle [a :: metric-space \Rightarrow b :: metric-space, real] \Rightarrow bool \rangle$   
**where**  $\langle lipschitz-with f \alpha \equiv lipschitz-with-on f \alpha UNIV \rangle$

**lemma** *lipschitz-with-onI* :  
 $\langle [0 \leq \alpha; \bigwedge x y. [x \in A; y \in A; x \neq y; f x \neq f y] \implies dist(f x) (f y) \leq \alpha * dist x y] \implies lipschitz-with-on f \alpha A \rangle$   
**unfolding** *lipschitz-with-on-def* **by** (*metis dist-eq-0-iff zero-le-dist zero-le-mult-iff*)

**lemma** *lipschitz-withI* :  
 $\langle [0 \leq \alpha; \bigwedge x y. x \neq y \implies f x \neq f y \implies dist(f x) (f y) \leq \alpha * dist x y] \implies lipschitz-with f \alpha \rangle$

by (rule lipschitz-with-onI)

**lemma** lipschitz-with-onD1 : ⟨lipschitz-with-on f α A ⟹ 0 ≤ α⟩  
**unfolding** lipschitz-with-on-def **by** simp

**lemma** lipschitz-withD1 : ⟨lipschitz-with f α ⟹ 0 ≤ α⟩  
**by** (rule lipschitz-with-onD1)

**lemma** lipschitz-with-onD2 :  
⟨lipschitz-with-on f α A ⟹ x ∈ A ⟹ y ∈ A ⟹ dist (f x) (f y) ≤  
 $\alpha * \text{dist } x y$   
**unfolding** lipschitz-with-on-def **by** simp

**lemma** lipschitz-withD2 :  
⟨lipschitz-with f α ⟹ dist (f x) (f y) ≤ α \* dist x y  
**unfolding** lipschitz-with-on-def **by** simp

**lemma** lipschitz-with-imp-lipschitz-with-on : ⟨lipschitz-with f α ⟹ lip-  
schitz-with-on f α A⟩  
**by** (simp add: lipschitz-with-on-def)

**lemma** lipschitz-with-on-imp-lipschitz-with-on-ge : ⟨lipschitz-with-on f  
 $\beta A$ ⟩  
**if** ⟨ $\alpha \leq \beta$ ⟩ **and** ⟨lipschitz-with-on f α A⟩  
**proof** (rule lipschitz-with-onI)  
**show** ⟨ $0 \leq \beta$ ⟩ **by** (meson order-trans lipschitz-with-onD1 that)  
**next**  
**fix** x y **assume** ⟨ $x \in A$ ⟩ **and** ⟨ $y \in A$ ⟩  
**from** ⟨lipschitz-with-on f α A⟩[THEN lipschitz-with-onD2, OF this]  
**have** ⟨ $\text{dist } (f x) (f y) \leq \alpha * \text{dist } x y$ ⟩ .  
**from** order-trans[OF this] **show** ⟨ $\text{dist } (f x) (f y) \leq \beta * \text{dist } x y$ ⟩  
**by** (simp add: mult-right-mono ⟨ $\alpha \leq \beta$ ⟩)  
**qed**

**theorem** lipschitz-with-on-comp-lipschitz-with-on :  
⟨lipschitz-with-on (λx. g (f x)) ( $\beta * \alpha$ ) A⟩  
**if** ⟨ $f`A \subseteq B$ ⟩ ⟨lipschitz-with-on g β B⟩ ⟨lipschitz-with-on f α A⟩  
**proof** (rule lipschitz-with-onI)  
**show** ⟨ $0 \leq \beta * \alpha$ ⟩ **by** (metis lipschitz-with-onD1 mult-nonneg-nonneg  
that(2, 3))  
**next**  
**fix** x y **assume** ⟨ $x \in A$ ⟩ **and** ⟨ $y \in A$ ⟩  
**with** ⟨ $f`A \subseteq B$ ⟩ **have** ⟨ $f x \in B$ ⟩ **and** ⟨ $f y \in B$ ⟩ **by** auto  
**have** ⟨ $\text{dist } (g (f x)) (g (f y)) \leq \beta * \text{dist } (f x) (f y)$ ⟩  
**by** (fact that(2)[THEN lipschitz-with-onD2, OF ⟨ $f x \in B$ ⟩ ⟨ $f y \in B$ ⟩])  
**also have** ⟨... ≤  $\beta * (\alpha * \text{dist } x y)$ ⟩

```

by (fact mult-left-mono[OF that(3)[THEN lipschitz-with-onD2, OF  

\x \in A \y \in A] that(2)[THEN lipschitz-with-onD1]])  

also have \dots = \beta * \alpha * dist x y by simp  

finally show dist (g (f x)) (g (f y)) \leq \dots .  

qed

```

```

corollary lipschitz-with-comp-lipschitz-with :  

\[lipschitz-with g \beta; lipschitz-with f \alpha\] \implies  

lipschitz-with (\lambda x. g (f x)) (\beta * \alpha)  

using lipschitz-with-on-comp-lipschitz-with-on by blast

```

### 1.1.2 Non-expanding Map

```

definition non-expanding-on :: \['a :: metric-space \implies 'b :: metric-space,  

'a set] \implies bool  

where non-expanding-on f A \equiv lipschitz-with-on f 1 A

```

```

abbreviation non-expanding :: \['a :: metric-space \implies 'b :: metric-space]  

\implies bool  

where non-expanding f \equiv non-expanding-on f UNIV

```

```

lemma non-expanding-onI :  

\[\forall x y. \[x \in A; y \in A; x \neq y; f x \neq f y\] \implies dist (f x) (f y) \leq dist  

x y\] \implies  

non-expanding-on f A  

by (simp add: lipschitz-with-onI non-expanding-on-def)

```

```

lemma non-expandingI :  

\[\forall x y. x \neq y \implies f x \neq f y \implies dist (f x) (f y) \leq dist x y\] \implies  

non-expanding f  

by (rule non-expanding-onI)

```

```

lemma non-expanding-onD :  

\non-expanding-on f A \implies x \in A \implies y \in A \implies dist (f x) (f y) \leq  

dist x y  

by (metis lipschitz-with-onD2 mult-1 non-expanding-on-def)

```

```

lemma non-expandingD : \non-expanding f \implies dist (f x) (f y) \leq dist  

x y  

by (simp add: non-expanding-onD)

```

```

lemma non-expanding-imp-non-expanding-on : \non-expanding f \implies  

non-expanding-on f A  

by (meson non-expandingD non-expanding-onI)

```

```

lemma non-expanding-on-comp-non-expanding-on :  

\[f ` A \subseteq B; non-expanding-on g B; non-expanding-on f A\] \implies

```

$\text{non-expanding-on } (\lambda x. g (f x)) A$   
**unfolding**  $\text{non-expanding-on-def}$   
**by** (*metis (no-types) lipschitz-with-on-comp-lipschitz-with-on mult-1*)

**corollary**  $\text{non-expanding-comp-non-expanding} :$   
 $\langle \llbracket \text{non-expanding } g; \text{ non-expanding } f \rrbracket \implies \text{non-expanding } (\lambda x. g (f x)) \rangle$   
**by** (*blast intro: non-expanding-on-comp-non-expanding-on*)

### 1.1.3 Contraction Map

**definition**  $\text{contraction-with-on} :: \langle [a :: \text{metric-space} \Rightarrow b :: \text{metric-space, real}, 'a \text{ set}] \Rightarrow \text{bool} \rangle$   
**where**  $\langle \text{contraction-with-on } f \alpha A \equiv \alpha < 1 \wedge \text{lipschitz-with-on } f \alpha A \rangle$

**abbreviation**  $\text{contraction-with} :: \langle [a :: \text{metric-space} \Rightarrow b :: \text{metric-space, real}] \Rightarrow \text{bool} \rangle$   
**where**  $\langle \text{contraction-with } f \alpha \equiv \text{contraction-with-on } f \alpha \text{ UNIV} \rangle$

**definition**  $\text{contraction-on} :: \langle [a :: \text{metric-space} \Rightarrow b :: \text{metric-space, 'a set}] \Rightarrow \text{bool} \rangle$   
**where**  $\langle \text{contraction-on } f A \equiv \exists \alpha. \text{contraction-with-on } f \alpha A \rangle$

**abbreviation**  $\text{contraction} :: \langle [a :: \text{metric-space} \Rightarrow b :: \text{metric-space}] \Rightarrow \text{bool} \rangle$   
**where**  $\langle \text{contraction } f \equiv \text{contraction-on } f \text{ UNIV} \rangle$

**lemma**  $\text{contraction-with-onI} :$   
 $\langle \llbracket 0 \leq \alpha; \alpha < 1; \bigwedge x y. \llbracket x \in A; y \in A; x \neq y; f x \neq f y \rrbracket \implies \text{dist} (f x) (f y) \leq \alpha * \text{dist} x y \rrbracket \implies \text{contraction-with-on } f \alpha A$   
**by** (*simp add: contraction-with-on-def lipschitz-with-onI*)

**lemma**  $\text{contraction-withI} :$   
 $\langle \llbracket 0 \leq \alpha; \alpha < 1; \bigwedge x y. x \neq y \implies f x \neq f y \implies \text{dist} (f x) (f y) \leq \alpha * \text{dist} x y \rrbracket \implies \text{contraction-with } f \alpha \rangle$   
**by** (*rule contraction-with-onI*)

**lemma**  $\text{contraction-with-onD1} : \langle \text{contraction-with-on } f \alpha A \implies 0 \leq \alpha \rangle$   
**by** (*simp add: contraction-with-on-def lipschitz-with-on-def*)

**lemma**  $\text{contraction-withD1} : \langle \text{contraction-with } f \alpha \implies 0 \leq \alpha \rangle$   
**by** (*simp add: contraction-with-onD1*)

**lemma**  $\text{contraction-with-onD2} : \langle \text{contraction-with-on } f \alpha A \implies \alpha <$

```

1>
by (simp add: contraction-with-on-def lipschitz-with-on-def)

lemma contraction-withD2 : <contraction-with f α ==> α < 1>
by (simp add: contraction-with-onD2)

lemma contraction-with-onD3 :
<contraction-with-on f α A ==> x ∈ A ==> y ∈ A ==> dist (f x) (f
y) ≤ α * dist x y
by (simp add: contraction-with-on-def lipschitz-with-on-def)

lemma contraction-withD3 : <contraction-with f α ==> dist (f x) (f y)
≤ α * dist x y
by (simp add: contraction-with-on-def lipschitz-withD2)

lemma contraction-with-imp-contraction-with-on:
<contraction-with f α ==> contraction-with-on f α A>
by (simp add: contraction-with-on-def lipschitz-with-imp-lipschitz-with-on)

lemma contraction-imp-contraction-on: <contraction f ==> contrac-
tion-on f A>
using contraction-on-def contraction-with-imp-contraction-with-on
by blast

lemma contraction-with-on-imp-contraction-on :
<contraction-with-on f α A ==> contraction-on f A>
unfolding contraction-on-def by blast

lemma contraction-with-imp-contraction: <contraction-with f α ==>
contraction f>
by (simp add: contraction-with-on-imp-contraction-on)

lemma contraction-onE:
<[contraction-on f A; ∧α. contraction-with-on f α A ==> thesis] ==>
thesis>
unfolding contraction-on-def by blast

lemma contractionE:
<[contraction f; ∧α. contraction-with f α ==> thesis] ==> thesis>
by (elim contraction-onE)

lemma contraction-with-on-imp-contraction-with-on-ge :
<[α ≤ β; β < 1; contraction-with-on f α A] ==> contraction-with-on
f β A>
by (simp add: contraction-with-on-def lipschitz-with-on-imp-lipschitz-with-on-ge)

```

## 1.2 Properties

```

lemma contraction-with-on-imp-lipschitz-with-on[simp] :
  <contraction-with-on f α A ==> lipschitz-with-on f α A>
  by (simp add: contraction-with-on-def)

lemma non-expanding-on-imp-lipschitz-with-one-on[simp] :
  <non-expanding-on f A ==> lipschitz-with-on f 1 A>
  by (simp add: non-expanding-on-def)

lemma contraction-on-imp-non-expanding-on[simp] :
  <contraction-on f A ==> non-expanding-on f A>
proof (elim contraction-onE, rule non-expanding-onI)
  fix α x y assume <x ∈ A> and <y ∈ A> and contra : <contraction-with-on f α A>
  show <dist (f x) (f y) ≤ dist x y>
    by (rule order-trans[OF contra[THEN contraction-with-onD3, OF
      <x ∈ A> <y ∈ A>]])
    (metis contra contraction-with-on-def mult-less-cancel-right1 nle-le
     order-less-le zero-le-dist)
  qed

lemma contraction-with-on-comp-contraction-with-on :
  <contraction-with-on (λx. g (f x)) (β * α) A>
  if <f ` A ⊆ B> <contraction-with-on g β B> <contraction-with-on f α
  A>
proof (unfold contraction-with-on-def, intro conjI)
  from that(2, 3)[THEN contraction-with-onD2] that(2)[THEN con-
  traction-with-onD1]
  show <β * α < 1> by (metis dual-order.strict-trans2 mult-less-cancel-left1
   nle-le order-less-le)
next
  show <lipschitz-with-on (λx. g (f x)) (β * α) A>
    by (rule lipschitz-with-on-comp-lipschitz-with-on[OF <f ` A ⊆ B>])
    (simp-all add: that(2, 3))
qed

corollary contraction-with-comp-contraction-with :
  <[contraction-with g β; contraction-with f α] ==> contraction-with
  (λx. g (f x)) (β * α)>
  by (blast intro: contraction-with-on-comp-contraction-with-on)

corollary contraction-on-comp-contraction-on :
  <[f ` A ⊆ B; contraction-on g B; contraction-on f A] ==> contrac-
  tion-on (λx. g (f x)) A>
proof (elim contraction-onE)
  fix α β assume <f ` A ⊆ B> <contraction-with-on g β B> <contrac-
  tion-with-on f α A>
  from contraction-with-on-comp-contraction-with-on[OF this]

```

**show**  $\langle \text{contraction-on } (\lambda x. g(fx)) A \rangle$  **by** (*fact contraction-with-on-imp-contraction-on*)  
**qed**

**corollary** *contraction-comp-contraction* :  
 $\langle \llbracket \text{contraction } g; \text{contraction } f \rrbracket \implies \text{contraction } (\lambda x. g(fx)) \rangle$   
**by** (*blast intro: contraction-on-comp-contraction-on*)

**lemma** *contraction-with-on-comp-non-expanding-on* :  
 $\langle \text{contraction-with-on } (\lambda x. g(fx)) \beta A \rangle$   
**if**  $\langle f' A \subseteq B \rangle \langle \text{contraction-with-on } g \beta B \rangle \langle \text{non-expanding-on } f A \rangle$   
**proof** (*unfold contraction-with-on-def, intro conjI*)  
**from** *that(2)[THEN contraction-with-onD2]* **show**  $\langle \beta < 1 \rangle$ .  
**next**  
**show**  $\langle \text{lipschitz-with-on } (\lambda x. g(fx)) \beta A \rangle$   
**by** (*rule lipschitz-with-on-comp-lipschitz-with-on[OF f' A ⊆ B, of g β 1, simplified]*)  
*(simp-all add: that(2, 3))*  
**qed**

**corollary** *contraction-with-comp-non-expanding* :  
 $\langle \llbracket \text{contraction-with } g \beta; \text{non-expanding } f \rrbracket \implies \text{contraction-with } (\lambda x. g(fx)) \beta \rangle$   
**by** (*blast intro: contraction-with-on-comp-non-expanding-on*)

**corollary** *contraction-on-comp-non-expanding-on* :  
 $\langle \llbracket f' A \subseteq B; \text{contraction-on } g B; \text{non-expanding-on } f A \rrbracket \implies \text{contraction-on } (\lambda x. g(fx)) A \rangle$   
**by** (*metis contraction-on-def contraction-with-on-comp-non-expanding-on*)

**corollary** *contraction-comp-non-expanding* :  
 $\langle \llbracket \text{contraction } g; \text{non-expanding } f \rrbracket \implies \text{contraction } (\lambda x. g(fx)) \rangle$   
**by** (*blast intro: contraction-on-comp-non-expanding-on*)

**lemma** *non-expanding-on-comp-contraction-with-on* :  
 $\langle \text{contraction-with-on } (\lambda x. g(fx)) \alpha A \rangle$   
**if**  $\langle f' A \subseteq B \rangle \langle \text{non-expanding-on } g B \rangle \langle \text{contraction-with-on } f \alpha A \rangle$   
**proof** (*unfold contraction-with-on-def, intro conjI*)  
**from** *that(3)[THEN contraction-with-onD2]* **show**  $\langle \alpha < 1 \rangle$ .  
**next**  
**show**  $\langle \text{lipschitz-with-on } (\lambda x. g(fx)) \alpha A \rangle$   
**by** (*rule lipschitz-with-on-comp-lipschitz-with-on[OF f' A ⊆ B, of g α 1, simplified]*)  
*(simp-all add: that(2, 3))*  
**qed**

**corollary** *non-expanding-comp-contraction-with* :  
 $\langle \llbracket \text{non-expanding } g; \text{contraction-with } f \alpha \rrbracket \implies \text{contraction-with } (\lambda x. g(fx)) \alpha \rangle$

```

 $g(f(x)) \alpha$ 
by (blast intro: non-expanding-on-comp-contraction-with-on)

corollary non-expanding-on-comp-contraction-on :
  ‹[f ‘ A ⊆ B; non-expanding-on g B; contraction-on f A] ⟹ contraction-on (λx. g(f x)) A›
by (metis contraction-on-def non-expanding-on-comp-contraction-with-on)

corollary non-expanding-comp-contraction :
  ‹[non-expanding g; contraction f] ⟹ contraction (λx. g(f x))›
by (blast intro: non-expanding-on-comp-contraction-on)

```

### 1.3 Banach's fixed-point Theorems

We rewrite the Banach's fixed-point theorems with our new definition.

```

theorem Banach-fix-type : <contraction f ⟹ ∃!x. f x = x>
  for f :: ‹'a :: complete-space ⇒ 'a›
  by (elim contractionE)
    (metis banach-fix-type contraction-withD1 contraction-withD2 contraction-withD3)

```

```

theorem Banach-fix:
  ‹contraction-on f s ⟹ ∃!x. x ∈ s ∧ f x = x› if <complete s> <s ≠ {}> <f ‘ s ⊆ s›
proof (elim contraction-onE, intro banach-fix[OF <complete s> <s ≠ {}> - - <f ‘ s ⊆ s>] ballI)
  show <contraction-with-on f α s ⟹ 0 ≤ α> for α by (fact contraction-with-onD1)
next
  show <contraction-with-on f α s ⟹ α < 1> for α by (fact contraction-with-onD2)
next
  show <contraction-with-on f α s ⟹ x ∈ s ⟹ y ∈ s ⟹
    dist(f x) (f y) ≤ α * dist x y> for α x y by (fact contraction-with-onD3)
qed

```

## 2 Locales factorizing the proof Work

### 2.1 Preliminaries on strictly decreasing Sequences

```

abbreviation strict-decseq :: <(nat ⇒ 'a :: order) ⇒ bool>
  where <strict-decseq ≡ monotone (<) (λx y. y < x)>

```

```

lemma strict-decseq-def : <strict-decseq X ⟷ (∀ m n. m < n → X

```

```

 $n < X m)$ 
by (fact monotone-def)

lemma strict-decseqI : <strict-decseq X> if < $\bigwedge n. X (\text{Suc } n) < X n$ >
by (metis Suc-le-eq decseqD decseq-Suc-iff le-less-trans nless-le strict-decseq-def
that)

lemma strict-decseqD : <strict-decseq X  $\implies$   $m < n \implies X n < X m$ >
using strict-decseq-def by blast

lemma strict-decseq-def-bis : <strict-decseq X  $\longleftrightarrow$  ( $\forall m n. X n < X$ 
 $m \longleftrightarrow m < n$ )>
by (metis linorder-less-linear order-less-imp-not-less strict-decseq-def)

lemma strict-decseq-def-ter : <strict-decseq X  $\longleftrightarrow$  ( $\forall m n. X n \leq X$ 
 $m \longleftrightarrow m \leq n$ )>
unfolding strict-decseq-def-bis
by (rule iffI, metis le-simps(2) order-le-less, simp add: less-le-not-le)

lemma strict-decseq-imp-decseq : <strict-decseq  $\sigma \implies$  decseq  $\sigma$ >
by (simp add: monotone-on-def order-le-less)

lemma strict-decseq-SucI : <( $\bigwedge n. X (\text{Suc } n) < X n$ )  $\implies$  strict-decseq
X>
by (metis Suc-le-eq decseqD decseq-Suc-iff less-le-not-le order-less-le
strict-decseq-def)

lemma strict-decseq-SucD : <strict-decseq A  $\implies$  A (Suc i) < A i>
by (simp add: strict-decseq-def)

```

Classically, a restriction space is given the structure of a metric space by defining  $\text{dist } x y \equiv \text{Inf} \{(1 / 2)^n \mid n. x \downarrow n = y \downarrow n\}$ . This obviously also works if we replace  $1 / 2$  by any real  $\delta$  such that  $0 < \delta$  and  $\delta < 1$ . But more generally, this still works if we set  $\text{dist } x y \equiv \text{Inf} \{\sigma n \mid n. x \downarrow n = y \downarrow n\}$  where  $\sigma$  is a sequence of *real* verifying  $\forall n. 0 < \sigma n$  and  $\sigma \longrightarrow 0$ . As you would expect, the more structure you have, the more powerful theorems you get. We explore all these variants in the theory below.

## 2.2 The Construction with Locales

Our formalization will extend the class *metric-space*. But some proofs are redundant, especially when it comes to the product type. So first we will be working with locales, and interpret them with the classes.

```

locale NonDecseqRestrictionSpace = PreorderRestrictionSpace +
— Factorization of the proof work.
fixes M :: "('a set) and restriction- $\sigma$  :: 'nat  $\Rightarrow$  real ( $\langle \sigma_{\downarrow} \rangle$ )
and restriction-dist :: ' $a \Rightarrow a \Rightarrow real$  ( $\langle dist_{\downarrow} \rangle$ )
assumes restriction- $\sigma$ -tendsto-zero :  $\langle restriction-{\sigma} \longrightarrow 0 \rangle$ 
and zero-less-restriction- $\sigma$  [simp] :  $\langle 0 < restriction-{\sigma} n \rangle$ 
and dist-restriction-is :
 $dist_{\downarrow} x y = (INF n \in restriction-related-set x y. restriction-{\sigma} n)$ 
begin

lemma zero-le-restriction- $\sigma$  [simp] :  $\langle 0 \leq \sigma_{\downarrow} n \rangle$ 
by (simp add: order-less-imp-le)

lemma restriction- $\sigma$ -neq-zero [simp] :  $\langle \sigma_{\downarrow} n \neq 0 \rangle$ 
by (metis zero-less-restriction- $\sigma$  order-less-irrefl)

lemma bounded-range-restriction- $\sigma$ :  $\langle bounded (range \sigma_{\downarrow}) \rangle$ 
by (fact convergent-imp-bounded[OF restriction- $\sigma$ -tendsto-zero])

abbreviation restriction- $\sigma$ -related-set :: ' $a \Rightarrow a \Rightarrow real$  set'
where  $\langle restriction-{\sigma}-related-set x y \equiv \sigma_{\downarrow} ` restriction-related-set x y \rangle$ 

abbreviation restriction- $\sigma$ -not-related-set :: ' $a \Rightarrow a \Rightarrow real$  set'
where  $\langle restriction-{\sigma}-not-related-set x y \equiv \sigma_{\downarrow} ` restriction-not-related-set x y \rangle$ 

lemma nonempty-restriction- $\sigma$ -related-set :
 $\langle restriction-{\sigma}-related-set x y \neq \{ \} \rangle$  by simp

lemma restriction- $\sigma$ -related-set-Un-restriction- $\sigma$ -not-related-set :
 $\langle restriction-{\sigma}-related-set x y \cup restriction-{\sigma}-not-related-set x y = range \sigma_{\downarrow} \rangle$ 
by blast

lemma ⟨bdd-above (restriction- $\sigma$ -related-set x y)⟩
by (meson bdd-above.I2 bdd-above.unfold bounded-imp-bdd-above
bounded-range-restriction- $\sigma$  rangeI)

lemma ⟨bdd-above (restriction- $\sigma$ -not-related-set x y)⟩
by (meson bdd-above.E bdd-above.I2 bounded-imp-bdd-above
bounded-range-restriction- $\sigma$  range-eqI)

lemma bounded-restriction- $\sigma$ -related-set:  $\langle bounded (restriction-{\sigma}-related-set x y) \rangle$ 

```

**by** (meson bounded-range-restriction- $\sigma$  bounded-subset image-mono top-greatest)

**lemma** bounded-restriction- $\sigma$ -not-related-set:  $\langle \text{bounded} (\text{restriction-}\sigma\text{-not-related-set } x y) \rangle$   
**by** (meson bounded-range-restriction- $\sigma$  bounded-subset image-mono subset-UNIV)

**corollary** restriction-space-Inf-properties:

$\langle a \in \text{restriction-}\sigma\text{-related-set } x y \implies \text{dist}_\downarrow x y \leq a \rangle$   
 $\langle [\forall a. a \in \text{restriction-}\sigma\text{-related-set } x y \implies b \leq a] \implies b \leq \text{dist}_\downarrow x y \rangle$   
**unfolding** dist-restriction-is  
**by** (simp-all add: bounded-has-Inf(1) bounded-restriction- $\sigma$ -related-set cInf-greatest)

**lemma** restriction- $\sigma$ -related-set-alt :  
 $\langle \text{restriction-}\sigma\text{-related-set } x y = \{\sigma_\downarrow n \mid n. n \in \text{restriction-related-set } x y\} \rangle$   
**by** blast

**lemma** exists-less-restriction- $\sigma$  :  $\langle \exists n. m < n \wedge \sigma_\downarrow n < \sigma_\downarrow m \rangle$   
**proof** (rule ccontr)  
**assume**  $\neg (\exists n > m. \sigma_\downarrow n < \sigma_\downarrow m)$   
**hence**  $\forall n \geq m. \sigma_\downarrow m \leq \sigma_\downarrow n$   
**by** (metis linorder-not-le nle-le)  
**hence**  $\neg \sigma_\downarrow \longrightarrow 0$   
**by** (meson Lim-bounded2 linorder-not-le zero-less-restriction- $\sigma$ )  
**with** restriction- $\sigma$ -tendsto-zero **show** False **by** simp  
**qed**

**lemma**  $\langle \text{dist}_\downarrow x y = \text{Inf} (\text{restriction-}\sigma\text{-related-set } x y) \rangle$   
**by** (fact dist-restriction-is)

**lemma** not-related-imp-dist-restriction-is-some-restriction- $\sigma$  :  
 $\langle \exists n. \text{dist}_\downarrow x y = \sigma_\downarrow n \wedge (\forall m \leq n. x \downarrow m \lesssim y \downarrow m) \rangle$  **if**  $\neg x \lesssim y$   
**proof** –  
**have**  $\langle \text{finite} (\text{restriction-related-set } x y) \rangle$   
**by** (simp add: finite-restriction-related-set-iff  $\neg x \lesssim y$ )  
**have**  $\langle \text{Inf} (\text{restriction-}\sigma\text{-related-set } x y) \in \text{restriction-}\sigma\text{-related-set } x y \rangle$   
**by** (rule closed-contains-Inf[OF nonempty-restriction- $\sigma$ -related-set])  
(simp-all add:  $\langle \text{finite} (\text{restriction-related-set } x y) \rangle$  finite-imp-closed)  
**hence**  $\langle \text{dist}_\downarrow x y \in \text{restriction-}\sigma\text{-related-set } x y \rangle$   
**by** (fold dist-restriction-is)  
**with** restriction-related-le **obtain** n  
**where**  $\langle n \in \text{restriction-related-set } x y \rangle$   $\langle \text{dist}_\downarrow x y = \sigma_\downarrow n \rangle$

```

⟨∀ m≤n. x ↓ m ≈ y ↓ m⟩ by blast
with ⟨dist↓ x y ∈ restriction-σ-related-set x y⟩ show ?thesis by blast
qed

lemma not-related-imp-dist-restriction-le-some-restriction-σ :
  ⟨¬ x ≈ y ⟹ ∃ n. dist↓ x y ≤ σ↓ n ∧ (¬ x ↓ Suc n ≈ y ↓ Suc n) ∧
  (∀ m≤n. x ↓ m ≈ y ↓ m)⟩
  by (blast intro: restriction-space-Inf-properties
    dest: ex-not-restriction-related-optimized)

lemma restriction-dist-eq-0-iff-related : ⟨dist↓ x y = 0 ⟺ x ≈ y⟩
proof (rule iffI)
  show ⟨dist↓ x y = 0 ⟹ x ≈ y⟩
  by (erule contrapos-pp)
    (auto dest: not-related-imp-dist-restriction-is-some-restriction-σ)
next
  show ⟨dist↓ x y = 0⟩ if ⟨x ≈ y⟩
  proof (rule order-antisym)
    show ⟨0 ≤ dist↓ x y⟩ by (simp add: cINF-greatest dist-restriction-is)
  next
    define Δ where ⟨Δ n ≡ - σ↓ n⟩ for n
    have * : ⟨Δ n ≤ - dist↓ x y⟩ for n
      unfolding Δ-def using ⟨x ≈ y⟩ restriction-space-Inf-properties(1)
        by simp (metis UNIV-restriction-related-set-iff rangeI)
      from restriction-σ-tendsto-zero tendsto-minus-cancel-left
      have ⟨Δ ⟶ 0⟩ unfolding Δ-def by force
        from lim-le[OF convergentI[OF ⟨Δ ⟶ 0⟩] *] limI[OF ⟨Δ ⟶ 0⟩]
        show ⟨dist↓ x y ≤ 0⟩ by simp
      qed
    qed
  end
end

locale DecseqRestrictionSpace = NonDecseqRestrictionSpace +
  assumes decseq-restriction-σ : ⟨decseq σ↓⟩
begin

lemma dist-restriction-is-bis :
  ⟨dist↓ x y = (if x ≈ y then 0 else σ↓ (Sup (restriction-related-set x y))))⟩
  if ⟨x ∈ M⟩ and ⟨y ∈ M⟩
proof (split if-split, intro conjI impI)
  from ⟨x ∈ M⟩ and ⟨y ∈ M⟩ show ⟨x ≈ y ⟹ restriction-dist x y = 0⟩

```

```

by (simp add: restriction-dist-eq-0-iff-related)
next
  show ‹dist↓ x y = σ↓ (Sup (restriction-related-set x y))› if ‹¬ x ≈ y›
    proof (rule order-antisym)
      show ‹dist↓ x y ≤ σ↓ (Sup (restriction-related-set x y))›
        unfolding dist-restriction-is
        by (rule cINF-lower[OF - Sup-in-restriction-related-set[OF ‹¬ x ≈ y›]])
          (meson bdd-below.I2 zero-le-restriction-σ)
    next
      show ‹σ↓ (Sup (restriction-related-set x y)) ≤ dist↓ x y›
        proof (unfold dist-restriction-is, rule cINF-greatest)
          show ‹restriction-related-set x y ≠ {}› by (fact nonempty-restriction-related-set)
        next
          fix n assume ‹n ∈ restriction-related-set x y›
          hence ‹n ≤ Sup (restriction-related-set x y)›
            by (metis ‹n ∈ restriction-related-set x y› le-cSup-finite
                finite-restriction-related-set-iff ‹¬ x ≈ y›)
            from decseqD[OF decseq-restriction-σ this]
            show ‹restriction-σ (Sup (restriction-related-set x y)) ≤ restriction-σ n› .
        qed
      qed
    qed

```

**lemma** *not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq* :

```

  ‹dist↓ x y = σ↓ (Sup (restriction-related-set x y))›
  ‹∀ m≤Sup (restriction-related-set x y). x↓ m ≈ y↓ m›
  ‹∀ m>Sup (restriction-related-set x y). ¬ x↓ m ≈ y↓ m›
  if ‹¬ x ≈ y› and ‹x ∈ M› and ‹y ∈ M›
    subgoal by (subst dist-restriction-is-bis; simp add: ‹¬ x ≈ y› ‹x ∈ M› ‹y ∈ M›)
    subgoal using Sup-in-restriction-related-set[OF ‹¬ x ≈ y›] restriction-related-le by blast
    using cSup-upper[OF - bdd-above-restriction-related-set-iff[THEN iffD2, OF ‹¬ x ≈ y›],
      of ‹Suc (Sup (restriction-related-set x y))›]
    by (metis (mono-tags, lifting) Suc-leI dual-order.refl mem-Collect-eq
        not-less-eq-eq restriction-related-le)

```

**theorem** *restriction-dist-tends-to-zero-independent-of-restriction-σ* :

- Very powerful theorem: the convergence of the distance to 0 is actually independent from the restriction sequence chosen.
- This is the theorem for which we had to work with locales first.

```

assumes <DecseqRestrictionSpace ( $\downarrow$ ) ( $\lesssim$ ) restriction- $\sigma'$  restriction-dist'>
  and < $\Sigma \in M$ > and <range  $\sigma \subseteq M$ >
shows <( $\lambda n.$  dist $\downarrow$  ( $\sigma$  n)  $\Sigma$ ) \longrightarrow 0 \longleftrightarrow ( $\lambda n.$  restriction-dist' ( $\sigma$  n)  $\Sigma$ ) \longrightarrow 0>
proof -
  { fix restriction- $\sigma$  restriction-dist restriction- $\sigma'$  restriction-dist'
    assume a1 : <DecseqRestrictionSpace ( $\downarrow$ ) ( $\lesssim$ ) restriction- $\sigma$  restriction-dist >
      and a2 : <DecseqRestrictionSpace ( $\downarrow$ ) ( $\lesssim$ ) restriction- $\sigma'$  restriction-dist'>
      and * : <( $\lambda n.$  restriction-dist ( $\sigma$  n)  $\Sigma$ ) \longrightarrow 0>

    interpret i1 : DecseqRestrictionSpace <( $\downarrow$ )> <( $\lesssim$ )> M restriction- $\sigma$ 
      restriction-dist by (fact a1)
    interpret i2 : DecseqRestrictionSpace <( $\downarrow$ )> <( $\lesssim$ )> M restriction- $\sigma'$ 
      restriction-dist' by (fact a2)

    have <( $\lambda n.$  restriction-dist' ( $\sigma$  n)  $\Sigma$ ) \longrightarrow 0>
    proof (rule metric-LIMSEQ-I)
      fix  $\varepsilon :: real$  assume < $0 < \varepsilon$ >

      from metric-LIMSEQ-D[OF i2.restriction- $\sigma$ -tends-to-zero < $0 < \varepsilon$ >]
      obtain N where ** : < $N \leq n \implies \text{restriction-}\sigma' n < \varepsilon$ > for n
      by auto

      fix N' :: nat

      have < $\exists N'. \forall n \geq N'. N \in \text{restriction-related-set} (\sigma n) \Sigma$ >
      proof (rule ccontr)
        assume < $\nexists N'. \forall n \geq N'. N \in \text{restriction-related-set} (\sigma n) \Sigma$ >
        hence *** : < $\forall N'. \exists n \geq N'. \neg \sigma n \downarrow N \lesssim \Sigma \downarrow N$ > by simp
          have **** : < $\forall N'. \exists n \geq N'. \text{restriction-}\sigma N \leq \text{restriction-dist}$ 
            ( $\sigma n$ )  $\Sigma$ >
            proof (rule allI)
              fix N' :: nat
              from *** obtain n where < $N' \leq n \wedge \neg \sigma n \downarrow N \lesssim \Sigma \downarrow N$ >
              by blast
              hence < $\neg \sigma n \lesssim \Sigma$ > using mono-restriction-related by blast
              have < $\text{restriction-}\sigma N \leq \text{restriction-dist} (\sigma n) \Sigma$ >
              proof (subst i1.dist-restriction-is-bis)
                show < $\sigma n \in M \wedge \Sigma \in M$ >
                  by (simp-all add: assms(2, 3) range-subsetD)
              next
              show < $\text{restriction-}\sigma N \leq (\text{if } \sigma n \lesssim \Sigma \text{ then } 0$ 
                else restriction- $\sigma (\text{Sup} (\text{restriction-related-set} (\sigma n) \Sigma)))$ >
                using < $\neg \sigma n \lesssim \Sigma \wedge \neg \sigma n \downarrow N \lesssim \Sigma \downarrow N$ > assms(2, 3)
  }

```

```

nle-le
  not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq(2)
    by (fastforce intro: decseqD[OF i1.decseq-restriction-σ])
  qed
  with ⟨N' ≤ n⟩ show ⟨∃ n ≥ N'. restriction-σ N ≤ restriction-dist
  (σ n) Σ⟩ by blast
  qed
  from metric-LIMSEQ-D[OF *, of ⟨restriction-σ N⟩]
  have ⟨∃ N''. ∀ n ≥ N''. restriction-dist (σ n) Σ < restriction-σ
  N⟩
    by (metis abs-of-nonneg i1.zero-le-restriction-σ i1.zero-less-restriction-σ
norm-conv-dist order-trans real-norm-def
      verit-comp-simplify1(3))
  with **** show False by fastforce
  qed

  then obtain N' where *** : ⟨N' ≤ n ⟹ N ∈ restriction-related-set (σ n) Σ⟩ for n by blast
  have ⟨restriction-dist' (σ n) Σ < ε⟩ if ⟨N' ≤ n⟩ for n
  proof (rule le-less-trans[OF i2.restriction-space-Inf-properties(1),
of ⟨restriction-σ' N⟩])
    from ⟨N' ≤ n⟩ ***
    show ⟨restriction-σ' N ∈ i2.restriction-σ-related-set (σ n) Σ⟩
  by blast
  next
    show ⟨restriction-σ' N < ε⟩ by (simp add: **)
  qed
  thus ⟨∃ N. ∀ n ≥ N. dist (restriction-dist' (σ n) Σ) 0 < ε⟩
  by (metis abs-of-nonneg dist-0-norm dist-commute i2.not-related-imp-dist-restriction-is-some-restriction-
i2.restriction-dist-eq-0-iff-related i2.zero-less-restriction-σ
order-less-imp-not-less real-norm-def
      verit-comp-simplify1(3))
  qed }
  note * = this

  show ⟨(λn. dist↓ (σ n) Σ) ⟶ 0 = (λn. restriction-dist' (σ n)
Σ) ⟶ 0⟩
  using * DecseqRestrictionSpace-axioms assms(1) by blast
qed

end

```

### 3 Ultrametric Structure of restriction Spaces

This has only be proven with the sort constraint, not inside the context of the class *metric-space* ...

```

context metric-space begin

lemma LIMSEQ-def : < $X \longrightarrow L \longleftrightarrow (\forall r>0. \exists no. \forall n\geq no. dist(X n) L < r)$ >
  unfolding tendsto-iff eventually-sequentially ..

lemma LIMSEQ-iff-nz: < $X \longrightarrow L \longleftrightarrow (\forall r>0. \exists no>0. \forall n\geq no. dist(X n) L < r)$ >
  by (meson Suc-leD LIMSEQ-def zero-less-Suc)

lemma metric-LIMSEQ-I: <( $\bigwedge r. 0 < r \implies \exists no. \forall n\geq no. dist(X n) L < r \implies X \longrightarrow L$ )>
  by (simp add: LIMSEQ-def)

lemma metric-LIMSEQ-D: < $X \longrightarrow L \implies 0 < r \implies \exists no. \forall n\geq no. dist(X n) L < r$ >
  by (simp add: LIMSEQ-def)

lemma LIMSEQ-dist-iff:
  < $f \longrightarrow l \longleftrightarrow (\lambda x. dist(f x) l) \longrightarrow 0$ >
proof (unfold LIMSEQ-def, rule iffI)
  show <( $\lambda n. dist(f n) l) \longrightarrow 0 \implies (\forall r>0. \exists no. \forall n\geq no. dist(f n) l < r)$ >
    by (metis (mono-tags, lifting) eventually-at-top-linorder order-tendstoD(2))
next
  show < $\forall r>0. \exists no. \forall n\geq no. dist(f n) l < r \implies (\lambda n. dist(f n) l) \longrightarrow 0$ >
    by (simp add: metric-space-class.LIMSEQ-def)
qed

lemma Cauchy-converges-subseq:
  fixes u::nat  $\Rightarrow$  'a
  assumes Cauchy u
  strict-mono r
   $(u \circ r) \longrightarrow l$ 
  shows u  $\longrightarrow l$ 
proof -
  have *: eventually ( $\lambda n. dist(u n) l < e$ ) sequentially if  $e > 0$  for e
  proof -
    have  $e/2 > 0$  using that by auto
    then obtain N1 where N1:  $\bigwedge m n. m \geq N1 \implies n \geq N1 \implies dist(u m) (u n) < e/2$ 
      using <Cauchy u> unfolding Cauchy-def by blast
    obtain N2 where N2:  $\bigwedge n. n \geq N2 \implies dist((u \circ r) n) l < e / 2$ 
      using order-tendstoD(2)[OF iffD1[OF LIMSEQ-dist-iff <( $u \circ r$ )  $\longrightarrow l$ ] <math>e/2 > 0</math>]

```

```

unfolding eventually-sequentially by auto
have dist (u n) l < e if n ≥ max N1 N2 for n
proof –
  have dist (u n) l ≤ dist (u n) ((u o r) n) + dist ((u o r) n) l
    by (rule dist-triangle)
  also have ... < e/2 + e/2
  proof (intro add-strict-mono)
    show dist (u n) ((u o r) n) < e / 2
    using N1[of n r n] N2[of n] that unfolding comp-def
    by (meson assms(2) le-trans max.bounded-iff strict-mono-imp-increasing)
    show dist ((u o r) n) l < e / 2
      using N2 that by auto
  qed
  finally show ?thesis by simp
  qed
  then show ?thesis unfolding eventually-sequentially by blast
  qed
have (λn. dist (u n) l) —→ 0
  by (simp add: less-le-trans * order-tendstoI)
then show ?thesis using LIMSEQ-dist-iff by auto
qed

end

```

### 3.1 The Construction with Classes

```

class restriction-σ = restriction-space +
  fixes restriction-σ :: «'a itself ⇒ nat ⇒ real» (⟨σ↓⟩)

```

```

setup ⟨Sign.add-const-constraint (const-name dist), NONE⟩
— To be able to use dist out of the metric-space class.

```

```

class non-decseq-restriction-space =
  uniformity-dist + open-uniformity + restriction-σ +
  — We do not assume the restriction sequence to be decseq yet.
  assumes restriction-σ-tendsto-zero' : ⟨σ↓ TYPE('a) —→ 0⟩
    and zero-less-restriction-σ' [simp] : ⟨0 < σ↓ TYPE('a) n⟩
    and dist-restriction-is' : ⟨dist x y = (INF n ∈ {n. x ↓ n = y ↓ n}. σ↓ TYPE('a) n)⟩
begin

```

```

sublocale NonDecseqRestrictionSpace ⟨(↓)⟩ ⟨(=)⟩ ⟨UNIV⟩ ⟨σ↓ TYPE('a)⟩ dist
  by unfold-locales (simp-all add: restriction-σ-tendsto-zero' dist-restriction-is')

```

```

end

setup ‹Sign.add-const-constraint (const-name ‹dist›, SOME typ ‹'a
:: metric-space  $\Rightarrow$  'a  $\Rightarrow$  real›)›
— Only allow dist in class metric-space (back to normal).

```

We hide duplicated facts  $\sigma_{\downarrow} \text{TYPE}(\text{'a}) \longrightarrow 0$   
 $0 < \sigma_{\downarrow} \text{TYPE}(\text{'a}) ?n$   
 $\text{dist} ?x ?y = \text{Inf} (\text{restriction-}\sigma\text{-related-set} ?x ?y)$ .

**hide-fact** restriction- $\sigma$ -tends-to-zero' zero-less-restriction- $\sigma$ ' dist-restriction-is'

**context** non-decseq-restriction-space **begin**

**subclass** ultrametric-space

**proof** unfold-locales

**show** ‹ $\text{dist} x y = 0 \longleftrightarrow x = y$ › **for** x y  
   **by** (simp add: restriction-dist-eq-0-iff-related)

**have** dist-commute : ‹ $\text{dist} x y = \text{dist} y x$ › **for** x y  
   **by** (simp add: dist-restriction-is) metis

**show** ‹ $\text{dist} x y \leq \max(\text{dist} x z, \text{dist} y z)$ › **for** x y z  
   **proof** –

**consider** ‹ $x \neq y$ › **and** ‹ $y \neq z$ › **and** ‹ $x \neq z$ › | ‹ $x = y \vee y = z \vee x = z$ › **by** blast

**thus** ‹ $\text{dist} x y \leq \max(\text{dist} x z, \text{dist} y z)$ ›

**proof** cases

**assume** ‹ $x \neq y$ › **and** ‹ $y \neq z$ › **and** ‹ $x \neq z$ ›

**from** this(1)[THEN not-related-imp-dist-restriction-le-some-restriction- $\sigma$ ]

      this(2, 3)[THEN not-related-imp-dist-restriction-is-some-restriction- $\sigma$ ]

**obtain** l m n

**where** \* : ‹ $\text{dist} x y \leq \sigma_{\downarrow} \text{TYPE}('a) l$ › ‹ $x \downarrow \text{Suc } l \neq y \downarrow \text{Suc } l$ ›

**and** \*\* : ‹ $\text{dist} y z = \sigma_{\downarrow} \text{TYPE}('a) m$ › ‹ $\forall k \leq m. y \downarrow k = z \downarrow k$ ›

**and** \*\*\* : ‹ $\text{dist} x z = \sigma_{\downarrow} \text{TYPE}('a) n$ › ‹ $\forall k \leq n. x \downarrow k = z \downarrow k$ ›

**by** blast

**have** ‹ $\min n m \leq l$ ›

**proof** (rule ccontr)

**assume** ‹ $\neg \min n m \leq l$ ›

**hence** ‹ $\text{Suc } l \leq n$ › **and** ‹ $\text{Suc } l \leq m$ › **by** simp-all

**with** \*\*(2) \*\*\*(2) Suc-le-lessD

**have** ‹ $x \downarrow \text{Suc } l = z \downarrow \text{Suc } l$ › ‹ $y \downarrow \text{Suc } l = z \downarrow \text{Suc } l$ › **by** simp-all

**hence** ‹ $x \downarrow \text{Suc } l = y \downarrow \text{Suc } l$ › **by** simp

**with** ‹ $x \downarrow \text{Suc } l \neq y \downarrow \text{Suc } l$ › **show** False ..

**qed**

**have** ‹ $\text{dist} x y \leq \sigma_{\downarrow} \text{TYPE}('a) (\min n m)$ ›

**by** (rule restriction-space-Inf-properties(1))

```

(simp add: **(2) ***(2))
also have ‹... ≤ max (dist x z) (dist y z)›
  unfolding **(1) ***(1) by linarith
  finally show ‹dist x y ≤ max (dist x z) (dist y z)› .
next
  show ‹x = y ∨ y = z ∨ x = z ⇒ dist x y ≤ max (dist x z) (dist y z)›
    by (elim disjE, simp-all add: dist-commute)
      (metis not-related-imp-dist-restriction-is-some-restriction-σ
       restriction-dist-eq-0-iff-related zero-le-restriction-σ dual-order.refl)
qed
qed
qed

end

context non-decseq-restriction-space begin

lemma restriction-tendsto-self: ‹(λn. x ↓ n) —→ x›
proof –
  have ‹(λn. dist (x ↓ n) x) —→ 0›
  proof (rule real-tendsto-sandwich[OF _])
    show ‹∀ F n in sequentially. 0 ≤ dist (x ↓ n) x› by simp
  next
    show ‹∀ F n in sequentially. dist (x ↓ n) x ≤ σ↓ TYPE('a) n›
      by (auto intro: eventually-sequentiallyI[OF restriction-space-Inf-properties(1)])
  next
    show ‹(λn. 0) —→ 0› by simp
  qed
  thus ‹(λn. x ↓ n) —→ x›
    by (subst tendsto-iff[of ‹(λn. x ↓ n)›], subst eventually-sequentially)
      (simp add: LIMSEQ-iff)
qed

end

class decseq-restriction-space = non-decseq-restriction-space +
  assumes decseq-restriction-σ': ‹decseq (σ↓ TYPE('a))›
begin

sublocale DecseqRestrictionSpace ‹(↓)› ‹(=)› ‹UNIV :: 'a set›
  ‹restriction-σ TYPE('a)› dist

```

by unfold-locales (simp add: decseq-restriction- $\sigma'$ )

— Removing  $x \in M$  and  $y \in M$ .  
**lemmas** dist-restriction-is-bis-simplified = dist-restriction-is-bis[simplified]  
**and** not-eq-imp-dist-restriction-is-restriction- $\sigma$ -Sup-restriction-eq-simplified  
= not-eq-imp-dist-restriction-is-restriction- $\sigma$ -Sup-restriction-eq[simplified]

**end**

We hide duplicated fact decseq ( $\sigma \downarrow \text{TYPE}('a')$ ).

**hide-fact** decseq-restriction- $\sigma'$

**class** strict-decseq-restriction-space = non-decseq-restriction-space +  
**assumes** strict-decseq-restriction- $\sigma$  : <strict-decseq ( $\sigma \downarrow \text{TYPE}('a')$ )>  
**begin**

**subclass** decseq-restriction-space  
**by** unfold-locales  
(fact strict-decseq-imp-decseq[OF strict-decseq-restriction- $\sigma$ ])

**end**

Generic Properties

**lemma** (in metric-space) dist-sequences-tendsto-zero-imp-tendsto-iff :  
 $\langle(\lambda n. \text{dist}(\sigma n) (\psi n)) \longrightarrow 0 \Rightarrow \sigma \longrightarrow \Sigma \leftrightarrow \psi \longrightarrow \Sigma\rangle$   
**proof** (rule iffI)  
show  $\langle\psi \longrightarrow \Sigma\rangle$  if  $\langle\sigma \longrightarrow \Sigma\rangle$  and  $\langle(\lambda n. \text{dist}(\sigma n) (\psi n)) \longrightarrow 0\rangle$  for  $\sigma \psi$   
**proof** –  
from that have \* :  $\langle(\lambda n. \text{dist}(\sigma n) (\psi n) + \text{dist}(\sigma n) \Sigma) \longrightarrow 0\rangle$   
unfolding LIMSEQ-dist-iff using tendsto-add-zero by blast  
have  $\langle(\lambda n. \text{dist}(\psi n) \Sigma) \longrightarrow 0\rangle$   
by (rule real-tendsto-sandwich  
[of  $\langle\lambda n. 0\rangle$   $\langle\lambda n. \text{dist}(\psi n) \Sigma\rangle$  -  $\langle\lambda n. \text{dist}(\sigma n) (\psi n) + \text{dist}(\sigma n) \Sigma\rangle$ ])  
(simp-all add: \* dist-triangle3)  
with LIMSEQ-dist-iff show  $\langle\psi \longrightarrow \Sigma\rangle$  by blast  
qed  
thus  $\langle(\lambda n. \text{dist}(\sigma n) (\psi n)) \longrightarrow 0 \Rightarrow \psi \longrightarrow \Sigma \Rightarrow \sigma \longrightarrow \Sigma\rangle$   
by (simp add: dist-commute)  
qed

```

lemma (in non-decseq-restriction-space) restricted-sequence-tendsto-iff
:
  ⟨(λn. σ n ↓ n) —→ Σ ↔ σ —→ Σ⟩
proof –
  have ⟨(λn. dist (σ n ↓ n) (σ n)) —→ 0⟩
  proof (unfold metric-space-class.LIMSEQ-def, intro allI impI)
    fix ε :: real assume ⟨0 < ε⟩
    from restriction-σ-tendsto-zero[unfolded metric-space-class.LIMSEQ-def,
    rule-format, OF ⟨0 < ε⟩]
    obtain no where ⟨∀ n≥no. σ↓ TYPE('a) n < ε⟩
      by (auto simp add: dist-real-def)
    have ⟨no ≤ n ==> dist (dist (σ n ↓ n) (σ n)) 0 < ε⟩ for n
      by (simp, rule le-less-trans[OF restriction-space-Inf-properties(1)[of
      ⟨σ↓ TYPE('a) n⟩]])
      (simp, simp add: ⟨∀ n≥no. σ↓ TYPE('a) n < ε⟩)
    thus ⟨∃ no. ∀ n≥no. dist (dist (σ n ↓ n) (σ n)) 0 < ε⟩ by blast
  qed

  thus ⟨(λn. σ n ↓ n) —→ Σ ↔ σ —→ Σ⟩
    by (simp add: dist-sequences-tendsto-zero-imp-tendsto-iff)
  qed

```

```

lemma (in non-decseq-restriction-space) Cauchy-restriction-chain :
  ⟨Cauchy σ⟩ if ⟨chain↓ σ⟩
proof (rule metric-CauchyI)
  fix ε :: real assume ⟨0 < ε⟩
  from LIMSEQ-D[OF restriction-σ-tendsto-zero ⟨0 < ε⟩, simplified]
  obtain M where ⟨σ↓ TYPE('a) M < ε⟩ by blast
  moreover have ⟨M ≤ m ==> M ≤ n ==> dist (σ m) (σ n) ≤ σ↓
  TYPE('a) M⟩ for m n
    by (rule restriction-space-Inf-properties(1), simp add: image-iff)
    (metis restriction-chain-def-ter ⟨chain↓ σ⟩)
  ultimately have ⟨∀ m≥M. ∀ n≥M. dist (σ m) (σ n) < ε⟩
    by (meson dual-order.strict-trans2)
  thus ⟨∃ M. ∀ m≥M. ∀ n≥M. dist (σ m) (σ n) < ε⟩ ..
  qed

```

```

lemma (in non-decseq-restriction-space) restriction-tendsto-imp-tendsto
:
  ⟨σ —→ Σ⟩ if ⟨σ —→ Σ⟩
proof (rule metric-LIMSEQ-I)
  fix ε :: real assume ⟨0 < ε⟩
  with restriction-σ-tendsto-zero[THEN LIMSEQ-D]
  obtain n0 where ⟨∀ k≥n0. σ↓ TYPE('a) k < ε⟩ by auto

```

```

hence  $\langle \sigma \downarrow \text{TYPE}('a) n0 < \varepsilon \rangle$  by simp
from restriction-tendstoD[OF  $\langle \sigma \dashrightarrow \Sigma \rangle$ ]
obtain  $n1$  where  $\langle \forall k \geq n1. \Sigma \downarrow n0 = \sigma k \downarrow n0 \rangle ..$ 
hence  $\langle \forall k \geq n1. \text{dist}(\sigma k) \Sigma \leq \sigma \downarrow \text{TYPE}('a) n0 \rangle$ 
    by (simp add: restriction-space-Inf-properties(1))
with  $\langle \sigma \downarrow \text{TYPE}('a) n0 < \varepsilon \rangle$ 
have  $\langle \forall k \geq n1. \text{dist}(\sigma k) \Sigma < \varepsilon \rangle$  by (meson order-le-less-trans)
thus  $\langle \exists n1. \forall k \geq n1. \text{dist}(\sigma k) \Sigma < \varepsilon \rangle ..$ 
qed

```

In Decseq Restriction Space

```

context decseq-restriction-space begin

lemma le-dist-to-restriction-eqE :
  obtains  $k$  where  $\langle n \leq k \rangle \langle \bigwedge x y :: 'a. \text{dist} x y \leq \sigma \downarrow \text{TYPE}('a) k \rangle$ 
   $\implies x \downarrow n = y \downarrow n$ 
proof –
  have  $\langle \exists k \geq n. \forall x y :: 'a. \text{dist} x y \leq \sigma \downarrow \text{TYPE}('a) k \longrightarrow x \downarrow n = y \downarrow n \rangle$ 
  proof (rule ccontr)
    assume  $\neg (\exists k \geq n. \forall x y :: 'a. \text{dist} x y \leq \sigma \downarrow \text{TYPE}('a) k \longrightarrow x \downarrow n = y \downarrow n)$ 
    hence  $\langle \forall k \geq n. \exists x y :: 'a. \text{dist} x y \leq \sigma \downarrow \text{TYPE}('a) k \wedge x \downarrow n \neq y \downarrow n \rangle$  by simp
    then obtain  $X Y :: \langle \text{nat} \Rightarrow 'a \rangle$ 
    where  $* : \langle \forall k \geq n. \text{dist}(X k) (Y k) \leq \sigma \downarrow \text{TYPE}('a) k \rangle$ 
       $\langle \forall k \geq n. X k \downarrow n \neq Y k \downarrow n \rangle$  by metis
    moreover obtain  $n0$  where  $\langle \forall k \geq n0. \sigma \downarrow \text{TYPE}('a) k < \sigma \downarrow \text{TYPE}('a) n \rangle$ 
    by (metis LIMSEQ-D zero-less-restriction- $\sigma$  abs-of-nonneg diff-zero
      restriction- $\sigma$ -tendsto-zero real-norm-def zero-le-restriction- $\sigma$ )
    ultimately have  $\langle \forall k \geq n+n0. \text{dist}(X k) (Y k) < \sigma \downarrow \text{TYPE}('a) n \rangle$ 
    by (metis dual-order.strict-trans2 add-leE)
    moreover from  $*(\mathcal{Q})$  have  $\langle \forall k \geq n. \sigma \downarrow \text{TYPE}('a) n \leq \text{dist}(X k) (Y k) \rangle$ 
    by (simp add: dist-restriction-is-bis)
    (metis (full-types) decseq-restriction- $\sigma$  [THEN decseqD] linorder-linear
      not-eq-imp-dist-restriction-is-restriction- $\sigma$ -Sup-restriction-eq-simplified(2))
    ultimately show False by (metis le-add1 linorder-not-le order-refl)
qed
  thus  $\langle (\bigwedge k. \llbracket n \leq k; \bigwedge x y :: 'a. \text{dist} x y \leq \sigma \downarrow \text{TYPE}('a) k \implies x \downarrow n = y \downarrow n \rrbracket \implies \text{thesis}) \implies \text{thesis} \rangle$  by blast
qed

```

```

theorem tendsto-iff-restriction-tendsto :  $\langle \sigma \longrightarrow \Sigma \longleftrightarrow \sigma \dashrightarrow \Sigma \rangle$ 
proof (rule iffI)

```

```

show ⟨σ ↴ → Σ ⟩ ⟹ σ —→ Σ by (fact restriction-tendsto-imp-tendsto)
next
  show ⟨σ ↴ → Σ⟩ if ⟨σ —→ Σ⟩
  proof (rule restriction-tendstoI)
    fix n
    obtain n0 where ⟨dist x y ≤ σ↓ TYPE('a) n0 ⟹ x ↓ n = y ↓
n⟩ for x y :: 'a
      by (metis le-dist-to-restriction-eqE)
  moreover from metric-LIMSEQ-D[OF ⟨σ —→ Σ⟩ zero-less-restriction-σ]
    obtain n1 where ⟨∀ k≥n1. dist (σ k) Σ < σ↓ TYPE('a) n0⟩ ..
  ultimately have ⟨∀ k≥n1. Σ ↓ n = σ k ↓ n⟩ by (metis dual-order.order-iff-strict)
    thus ⟨∃ n1. ∀ k≥n1. Σ ↓ n = σ k ↓ n⟩ ..
  qed
qed

```

**corollary** convergent-iff-restriction-convergent : ⟨convergent σ ↔ convergent↓ σ⟩  
**by** (simp add: convergent-def restriction-convergent-def tendsto-iff-restriction-tendsto)

**theorem** complete-iff-restriction-complete :  
⟨(∀σ. Cauchy σ —→ convergent σ) ↔ (∀σ. chain↓ σ —→ convergent↓ σ)⟩

— The following result shows that we have not lost anything with our definitions of convergence, completeness, etc. specific to restriction spaces.

**proof** (intro iffI impI allI)

fix σ assume hyp : ⟨∀σ. Cauchy σ —→ convergent σ⟩ and ⟨chain↓ σ⟩

from Cauchy-restriction-chain ⟨chain↓ σ⟩ have ⟨Cauchy σ⟩ by blast  
hence ⟨convergent σ⟩ by (simp add: hyp)  
thus ⟨convergent↓ σ⟩ by (simp add: convergent-iff-restriction-convergent)

next

fix σ assume hyp : ⟨∀σ. chain↓ σ —→ convergent↓ σ⟩ and ⟨Cauchy σ⟩

from ⟨Cauchy σ⟩ have \* : ⟨∀ n. ∃ k. ∀ l≥k. σ l ↓ n = σ k ↓ n⟩  
by (metis (mono-tags, opaque-lifting) Cauchy-altdef2 dual-order.order-iff-strict  
le-dist-to-restriction-eqE zero-less-restriction-σ)

define f where ⟨f ≡ rec-nat

(LEAST k. ∀ l≥k. σ l ↓ 0 = σ k ↓ 0)  
(λn k. LEAST l. k < l ∧ (∀ m≥l. σ m ↓ Suc n = σ l  
↓ Suc n))⟩

have f-Suc-def : ⟨f (Suc n) = (LEAST l. f n < l ∧ (∀ m≥l. σ m ↓  
Suc n = σ l ↓ Suc n))⟩

(is ⟨f (Suc n) = Least (?f-Suc n)⟩) for n by (simp add: f-def)  
from \* have \*\* : ⟨∃ k>f n. ∀ m≥k. σ m ↓ Suc n = σ k ↓ Suc n⟩ for  
n  
by (metis dual-order.trans lessI linorder-not-le order-le-less)

```

have ⟨strict-mono f⟩
proof (unfold strict-mono-Suc-iff, rule allI)
  show ⟨f n < f (Suc n)⟩ for n
    by (fact LeastI-ex[of ⟨?f-Suc n⟩, folded f-Suc-def, OF **, THEN
conjunct1])
qed

have ⟨chain↓ (λn. (σ ∘ f) n ↓ n)⟩ — key point
proof (rule restriction-chainI)
  fix n
  have ⟨σ (f (Suc n)) ↓ n = σ (f n) ↓ n⟩
  proof (cases n)
    show ⟨n = 0 ⟹ σ (f (Suc n)) ↓ n = σ (f n) ↓ n⟩ by simp
  next
    show ⟨n = Suc nat ⟹ σ (f (Suc n)) ↓ n = σ (f n) ↓ n⟩ for nat
      proof (clarify, intro LeastI-ex[of ⟨?f-Suc nat⟩,
folded f-Suc-def, THEN conjunct2, rule-format, OF **])
        show ⟨n = Suc nat ⟹ f (Suc nat) ≤ f (Suc (Suc nat))⟩
          by (simp add: ⟨strict-mono f⟩ strict-mono-less-eq)
      qed
  qed
  thus ⟨(σ ∘ f) (Suc n) ↓ Suc n ↓ n = (σ ∘ f) n ↓ n⟩ by simp
qed

with hyp have ⟨convergent↓ (λn. (σ ∘ f) n ↓ n)⟩ by simp
hence ⟨convergent↓ (λn. (σ ∘ f) n)⟩
  by (simp add: restriction-convergent-restricted-iff-restriction-convergent)
hence ⟨convergent (λn. (σ ∘ f) n)⟩
  by (simp add: convergent-iff-restriction-convergent)
with Cauchy-converges-subseq[OF ⟨Cauchy σ⟩ ⟨strict-mono f⟩]
show ⟨convergent σ⟩ unfolding convergent-def by blast
qed

end

```

The following classes will be useful later.

```

class complete-decseq-restriction-space = decseq-restriction-space +
  assumes restriction-chain-imp-restriction-convergent' : ⟨chain↓ σ
  ⟹ convergent↓ σ⟩
begin

  subclass complete-restriction-space
    by unfold-locales (fact restriction-chain-imp-restriction-convergent')

  subclass complete-ultrametric-space
  proof (unfold-locales)
    from complete-iff-restriction-complete restriction-chain-imp-restriction-convergent
    show ⟨Cauchy σ ⟹ convergent σ⟩ for σ by blast
  qed

```

**end**

We hide duplicated fact  $\text{chain}_{\downarrow} ?\sigma \implies \text{convergent}_{\downarrow} ?\sigma$ .

**hide-fact** *restriction-chain-imp-restriction-convergent'*

```
class complete-strict-decseq-restriction-space = strict-decseq-restriction-space
+
assumes restriction-chain-imp-restriction-convergent' : <chain_{\downarrow} \sigma
implies convergent_{\downarrow} \sigma>
begin

subclass complete-decseq-restriction-space
by (unfold-locales) (fact restriction-chain-imp-restriction-convergent')

end
```

We hide duplicated fact  $\text{chain}_{\downarrow} ?\sigma \implies \text{convergent}_{\downarrow} ?\sigma$ .

**hide-fact** *restriction-chain-imp-restriction-convergent'*

```
class restriction-\delta = restriction-\sigma +
fixes restriction-\delta :: <'a itself \Rightarrow real> (<\delta_{\downarrow}>)
assumes zero-less-restriction-\delta [simp] : <0 < \delta_{\downarrow} \text{TYPE}('a)>
and restriction-\delta-less-one [simp] : <\delta_{\downarrow} \text{TYPE}('a) < 1>
begin

lemma zero-le-restriction-\delta [simp] : <0 \leq \delta_{\downarrow} \text{TYPE}('a)>
and restriction-\delta-le-one [simp] : <\delta_{\downarrow} \text{TYPE}('a) \leq 1>
and zero-le-pow-restriction-\delta [simp] : <0 \leq \delta_{\downarrow} \text{TYPE}('a) ^ n>
by (simp-all add: order-less-imp-le)

lemma pow-restriction-\delta-le-one [simp] : <\delta_{\downarrow} \text{TYPE}('a) ^ n \leq 1>
by (simp add: power-le-one)

lemma pow-restriction-\delta-less-one [simp] : <n \neq 0 \implies \delta_{\downarrow} \text{TYPE}('a) ^ n < 1>
by (metis restriction-\delta-less-one zero-less-restriction-\delta not-gr-zero
power-0 power-strict-decreasing)

end
```

**setup** <*Sign.add-const-constraint (const-name dist, NONE)*>  
— To be able to use *dist* out of the *metric-space* class.

```
class at-least-geometric-restriction-space =
uniformity-dist + open-uniformity + restriction-\delta +
```

```

assumes zero-less-restriction- $\sigma'$  :  $\langle 0 < \sigma_\downarrow \text{TYPE}('a) n \rangle$ 
and restriction- $\sigma$ -le :
 $\langle \sigma_\downarrow \text{TYPE}('a) (\text{Suc } n) \leq \delta_\downarrow \text{TYPE}('a) * \sigma_\downarrow \text{TYPE}('a) n \rangle$ 
and dist-restriction-is' :
 $\langle \text{dist } x y = (\text{INF } n \in \{n. x \downarrow n = y \downarrow n\}. \sigma_\downarrow \text{TYPE}('a) n) \rangle$ 

setup ⟨Sign.add-const-constraint (const-name dist, SOME typ ⟨'a
:: metric-space  $\Rightarrow$  'a  $\Rightarrow$  real⟩)
— Only allow dist in class metric-space (back to normal).

context at-least-geometric-restriction-space begin

lemma restriction- $\sigma$ -le-restriction- $\sigma$ -times-pow-restriction- $\delta$  :
 $\langle \sigma_\downarrow \text{TYPE}('a) (n + k) \leq \sigma_\downarrow \text{TYPE}('a) n * \delta_\downarrow \text{TYPE}('a) ^ k \rangle$ 
by (induct k, simp-all)
(metis dual-order.trans restriction- $\sigma$ -le zero-less-restriction- $\delta$ 
mult.left-commute mult-le-cancel-left-pos)

lemma restriction- $\sigma$ -le-pow-restriction- $\delta$  :
 $\langle \sigma_\downarrow \text{TYPE}('a) n \leq \sigma_\downarrow \text{TYPE}('a) 0 * \delta_\downarrow \text{TYPE}('a) ^ n \rangle$ 
by (metis add-0 restriction- $\sigma$ -le-restriction- $\sigma$ -times-pow-restriction- $\delta$ )

subclass strict-decseq-restriction-space
proof unfold-locales
have  $\langle \forall_F n \text{ in sequentially}. 0 \leq \sigma_\downarrow \text{TYPE}('a) n \rangle$ 
by (simp add: zero-less-restriction- $\sigma'$  order-less-imp-le)
moreover have  $\langle \forall_F n \text{ in sequentially}. \sigma_\downarrow \text{TYPE}('a) n$ 
 $\leq \sigma_\downarrow \text{TYPE}('a) 0 * \delta_\downarrow \text{TYPE}('a) ^ n \rangle$ 
by (simp add: restriction- $\sigma$ -le-pow-restriction- $\delta$ )
moreover have  $\langle (\lambda n. 0) \longrightarrow 0 \rangle$  by simp
moreover have  $\langle (\lambda n. \sigma_\downarrow \text{TYPE}('a) 0 * \delta_\downarrow \text{TYPE}('a) ^ n) \longrightarrow$ 
 $0 \rangle$ 
by (simp add: LIMSEQ-abs-realpow-zero2 abs-of-pos tends-to-mult-right-zero)
ultimately show  $\langle \sigma_\downarrow \text{TYPE}('a) \longrightarrow 0 \rangle$  by (fact real-tends-to-sandwich)
next
show  $\langle 0 < \sigma_\downarrow \text{TYPE}('a) n \rangle$  for n
by (fact zero-less-restriction- $\sigma'$ )
next
show  $\langle \text{dist } x y = \text{Inf } (\sigma_\downarrow \text{TYPE}('a) ` \{n. x \downarrow n = y \downarrow n\}) \rangle$  for x y
by (fact dist-restriction-is')
next
show  $\langle \text{strict-decseq } (\sigma_\downarrow \text{TYPE}('a)) \rangle$ 
by (rule strict-decseqI, rule le-less-trans[OF restriction- $\sigma$ -le])
(simp add: zero-less-restriction- $\sigma'$ )
qed

```

```
lemma < $0 < \delta_{\downarrow} \text{TYPE}'(a) \wedge n$ > by simp
```

```
end
```

We hide duplicated facts  $0 < \sigma_{\downarrow} \text{TYPE}'(a) ?n$   
 $\text{dist} ?x ?y = \text{Inf} (\text{restriction-}\sigma\text{-related-set} ?x ?y)$ .

```
hide-fact zero-less-restriction- $\sigma'$  dist-restriction-is'
```

```
class complete-at-least-geometric-restriction-space = at-least-geometric-restriction-space
+
assumes restriction-chain-imp-restriction-convergent' : < $\text{chain}_{\downarrow} \sigma$ 
 $\implies \text{convergent}_{\downarrow} \sigma$ >
begin
```

```
subclass complete-strict-decseq-restriction-space
by unfold-locales (fact restriction-chain-imp-restriction-convergent')
```

```
end
```

We hide duplicated fact  $\text{chain}_{\downarrow} ?\sigma \implies \text{convergent}_{\downarrow} ?\sigma$ .

```
hide-fact restriction-chain-imp-restriction-convergent'
```

```
setup < $\text{Sign.add-const-constraint} (\text{const-name} \langle \text{dist} \rangle, \text{NONE})$ >
— To be able to use  $\text{dist}$  out of the metric-space class.
```

```
class geometric-restriction-space = uniformity-dist + open-uniformity
+ restriction- $\delta$  +
assumes restriction- $\sigma$ -is : < $\sigma_{\downarrow} \text{TYPE}'(a) n = \delta_{\downarrow} \text{TYPE}'(a) \wedge n$ >
    and dist-restriction-is' : < $\text{dist} x y = (\text{INF } n \in \{n. x \downarrow n = y \downarrow n\}.$ 
 $\sigma_{\downarrow} \text{TYPE}'(a) n)$ >
begin
```

This is what “restriction space” usually mean in the literature.  
The previous classes are generalizations of this concept (even this one is a generalization, since we usually have  $\delta_{\downarrow} \text{TYPE}'(a) = 1 / 2$ ).

```
subclass at-least-geometric-restriction-space
proof unfold-locales
    show < $0 < \sigma_{\downarrow} \text{TYPE}'(a) n$ > for n by (simp add: restriction- $\sigma$ -is)
next
```

```

show  $\langle \sigma_{\downarrow} \text{TYPE('a)} (\text{Suc } n) \leq \delta_{\downarrow} \text{TYPE('a)} * \sigma_{\downarrow} \text{TYPE('a)} n \rangle$  for
 $n$ 
    by (simp add: restriction- $\sigma$ -is)
next
    show  $\langle \text{dist } x y = \text{Inf} (\sigma_{\downarrow} \text{TYPE('a)} ' \{n. x \downarrow n = y \downarrow n\}) \rangle$  for  $x y$ 
        by (fact dist-restriction-is')
qed

lemma  $\langle 0 < \delta_{\downarrow} \text{TYPE('a)} \wedge n \rangle$  by simp
end

setup  $\langle \text{Sign.add-const-constraint } (\text{const-name } \langle \text{dist} \rangle, \text{SOME typ } 'a :: \text{metric-space} \Rightarrow 'a \Rightarrow \text{real}) \rangle$ 
— Only allow dist in class metric-space (back to normal).

```

We hide duplicated fact *dist*  $?x ?y = \text{Inf}$  (*restriction- $\sigma$ -related-set*  $?x ?y$ ).

**hide-fact** *dist-restriction-is'*

```

class complete-geometric-restriction-space = geometric-restriction-space
+
assumes restriction-chain-imp-restriction-convergent' :  $\langle \text{chain}_{\downarrow} \sigma \Rightarrow \text{convergent}_{\downarrow} \sigma \rangle$ 
begin

subclass complete-at-least-geometric-restriction-space
    by (unfold-locales) (fact restriction-chain-imp-restriction-convergent')
end

```

We hide duplicated fact *chain* $_{\downarrow} \sigma \Rightarrow \text{convergent}_{\downarrow} \sigma$ .

**hide-fact** *restriction-chain-imp-restriction-convergent'*

```

theorem geometric-restriction-space-completeI :  $\langle \text{convergent } \sigma \rangle$ 
    if  $\langle \bigwedge \sigma :: \text{nat} \Rightarrow 'a. \text{restriction-chain } \sigma \Rightarrow \exists \Sigma. \forall n. \Sigma \downarrow n = \sigma n \rangle$ 
        and  $\langle \text{Cauchy } \sigma \rangle$  for  $\sigma :: \langle \text{nat} \Rightarrow 'a :: \text{geometric-restriction-space} \rangle$ 
        by (metis complete-iff-restriction-complete convergent-def
            convergent-iff-restriction-convergent ext restriction-tendsto-self
            that)

```

— Because *cball* is not defined in *metric-space*.

**lemma** (**in** *non-decseq-restriction-space*) *restriction-cball-subset-cball* :

$\langle \sigma \downarrow \text{TYPE}('a) n \leq r \implies \mathcal{B}_\downarrow(\Sigma, n) \subseteq \{x. \text{dist } \Sigma x \leq r\} \rangle$   
**by** (simp add: subset-iff restriction-cball-mem-iff)  
 (simp add: dual-order.trans restriction-space-Inf-properties(1))

**corollary** restriction-cball-subset-cball-bis :  
 $\langle \sigma \downarrow \text{TYPE}('a) n \leq r \implies \mathcal{B}_\downarrow(\Sigma, n) \subseteq \text{cball } \Sigma r \rangle$   
**for**  $\Sigma :: \langle 'a :: \text{non-decseq-restriction-space} \rangle$   
**unfolding** cball-def **by** (fact restriction-cball-subset-cball)

**lemma** (in non-decseq-restriction-space) restriction-ball-subset-ball :  
 $\langle \sigma \downarrow \text{TYPE}('a) n < r \implies \text{restriction-ball } \Sigma n \subseteq \{x. \text{dist } \Sigma x < r\} \rangle$   
**by** (simp add: subset-iff restriction-cball-mem-iff)  
 (metis (mono-tags, lifting) image-eqI mem-Collect-eq order-le-less-trans  
 restriction-related-pred restriction-space-Inf-properties(1))

**corollary** restriction-ball-subset-ball-bis :  
 $\langle \sigma \downarrow \text{TYPE}('a) n < r \implies \text{restriction-ball } \Sigma n \subseteq \text{ball } \Sigma r \rangle$   
**for**  $\Sigma :: \langle 'a :: \text{non-decseq-restriction-space} \rangle$   
**unfolding** ball-def **by** (fact restriction-ball-subset-ball)

**lemma** (in strict-decseq-restriction-space)  
 $\text{restriction-cball-is-cball} : \langle \mathcal{B}_\downarrow(\Sigma, n) = \{x. \text{dist } \Sigma x \leq \sigma \downarrow \text{TYPE}('a) n\} \rangle$   
**proof** (intro subset-antisym subsetI)  
**from** restriction-cball-subset-cball  
**show**  $\langle x \in \mathcal{B}_\downarrow(\Sigma, n) \implies x \in \{x. \text{dist } \Sigma x \leq \sigma \downarrow \text{TYPE}('a) n\} \rangle$  **for**  $x$  **by** blast  
**next**  
**show**  $\langle x \in \mathcal{B}_\downarrow(\Sigma, n) \rangle$  **if**  $\langle x \in \{x. \text{dist } \Sigma x \leq \sigma \downarrow \text{TYPE}('a) n\} \rangle$  **for**  $x$   
**proof** (cases  $\langle \Sigma = x \rangle$ )  
**show**  $\langle \Sigma = x \implies x \in \mathcal{B}_\downarrow(\Sigma, n) \rangle$  **by** simp  
**next**  
**assume**  $\langle \Sigma \neq x \rangle$   
**with** not-related-imp-dist-restriction-is-some-restriction- $\sigma$  **obtain**  
 $m$   
**where**  $\langle \text{dist } \Sigma x = \sigma \downarrow \text{TYPE}('a) m \rangle$   $\langle \forall k \leq m. \Sigma \downarrow k = x \downarrow k \rangle$  **by** blast  
**from**  $\langle x \in \{x. \text{dist } \Sigma x \leq \sigma \downarrow \text{TYPE}('a) n\} \rangle$   
**have**  $\langle \text{dist } \Sigma x \leq \sigma \downarrow \text{TYPE}('a) n \rangle$  **by** simp  
**with**  $\langle \text{dist } \Sigma x = \sigma \downarrow \text{TYPE}('a) m \rangle$  **have**  $\langle n \leq m \rangle$   
**by** (metis linorder-not-less strict-decseq-restriction- $\sigma$  strict-decseq-def-bis)  
**with**  $\langle \forall k \leq m. \Sigma \downarrow k = x \downarrow k \rangle$  **have**  $\langle \Sigma \downarrow n = x \downarrow n \rangle$  **by** simp  
**thus**  $\langle x \in \mathcal{B}_\downarrow(\Sigma, n) \rangle$  **by** (simp add: restriction-cball-mem-iff)  
**qed**  
**qed**

```

lemma restriction-cball-is-cball-bis : < $\mathcal{B}_\downarrow(\Sigma, n) = \text{cball } \Sigma (\sigma_\downarrow \text{TYPE}('a) n)$ >
  for  $\Sigma :: ('a :: \text{strict-decseq-restriction-space})$ 
  by (simp add: cball-def restriction-cball-is-cball)

lemma (in strict-decseq-restriction-space)
  restriction-ball-is-ball : < $\text{restriction-ball } \Sigma n = \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\}$ >
proof (intro subset-antisym subsetI)
  show < $x \in \text{restriction-ball } \Sigma n \implies x \in \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\}$ > for  $x$ 
    by (simp add: subset-iff)
    (metis lessI local.restriction-cball-is-cball strict-decseq-restriction-σ
      mem-Collect-eq order-le-less-trans strict-decseq-def-bis)
next
  show < $x \in \text{restriction-ball } \Sigma n \text{ if } x \in \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\}$ > for  $x$ 
    proof (cases < $\Sigma = x$ >)
      show < $\Sigma = x \implies x \in \text{restriction-ball } \Sigma n$ > by simp
next
  assume < $\Sigma \neq x$ >
  with not-related-imp-dist-restriction-is-some-restriction-σ obtain
   $m$ 
    where < $\text{dist } \Sigma x = \sigma_\downarrow \text{TYPE}('a) m \wedge \forall k \leq m. \Sigma \downarrow k = x \downarrow k$ > by blast
    from < $x \in \{x. \text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n\}$ >
    have < $\text{dist } \Sigma x < \sigma_\downarrow \text{TYPE}('a) n$ > by simp
    with < $\text{dist } \Sigma x = \sigma_\downarrow \text{TYPE}('a) m$ > have < $n < m$ >
      using strict-decseq-restriction-σ strict-decseq-def-bis by auto
    with < $\forall k \leq m. \Sigma \downarrow k = x \downarrow k$ > have < $\Sigma \downarrow \text{Suc } n = x \downarrow \text{Suc } n$ > by
      simp
    thus < $x \in \text{restriction-ball } \Sigma n$ > by (simp add: restriction-cball-mem-iff)
    qed
qed

lemma restriction-ball-is-ball-bis : < $\text{restriction-ball } \Sigma n = \text{ball } \Sigma (\sigma_\downarrow \text{TYPE}('a) n)$ >
  for  $\Sigma :: ('a :: \text{strict-decseq-restriction-space})$ 
  by (simp add: ball-def restriction-ball-is-ball)

```

```

lemma isCont-iff-restriction-cont-at : < $\text{isCont } f \Sigma \longleftrightarrow \text{restriction-cont-at } f \Sigma$ >
  for  $f :: ('a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{decseq-restriction-space})$ 
  by (unfold restriction-cont-at-def continuous-at-sequentially comp-def,

```

*fold tendsto-iff-restriction-tendsto) simp*

```

lemma (in strict-decseq-restriction-space)
  open-iff-restriction-open : <open U  $\longleftrightarrow$  open↓ U>
proof (unfold open-dist restriction-open-iff-restriction-cball-characterization,
intro iffI ballI)
  fix  $\Sigma$  assume < $\forall x \in U. \exists e > 0. \forall y. dist y x < e \longrightarrow y \in U$ > and < $\Sigma$ 
 $\in U$ >
  then obtain  $e$  where < $0 < e$ > < $\forall y. dist y \Sigma < e \longrightarrow y \in U$ > by
  blast
  from < $0 < e$ > obtain  $n$  where < $\sigma_{\downarrow} TYPE('a) n < e$ >
    by (metis eventually-at-top-linorder le-refl restriction-σ-tendsto-zero
order-tendstoD(2))
  with < $\forall y. dist y \Sigma < e \longrightarrow y \in U$ > dist-commute restriction-cball-is-cball
  have < $\mathcal{B}_{\downarrow}(\Sigma, n) \subseteq U$ > by auto
  thus < $\exists n. \mathcal{B}_{\downarrow}(\Sigma, n) \subseteq U$ > ..
next
  fix  $x$  assume < $\forall \Sigma \in U. \exists n. \mathcal{B}_{\downarrow}(\Sigma, n) \subseteq U$ > < $x \in U$ >
  then obtain  $n$  where < $\mathcal{B}_{\downarrow}(x, n) \subseteq U$ > by blast
  hence < $\forall y. dist y x < \sigma_{\downarrow} TYPE('a) n \longrightarrow y \in U$ >
    by (simp add: dist-commute restriction-cball-is-cball subset-iff)
  thus < $\exists e > 0. \forall y. dist y x < e \longrightarrow y \in U$ >
    using zero-less-restriction-σ' by blast
qed

```

```

lemma (in strict-decseq-restriction-space)
  closed-iff-restriction-closed : <closed U  $\longleftrightarrow$  closed↓ U>
by (simp add: closed-open open-iff-restriction-open restriction-open-Compl-iff)

```

```

lemma continuous-on-iff-restriction-cont-on :
  < $\text{open } U \implies \text{continuous-on } U f \longleftrightarrow \text{restriction-cont-on } f U$ >
  for  $f :: ('a :: \text{decseq-restriction-space} \Rightarrow 'b :: \text{decseq-restriction-space})$ 
  by (simp add: restriction-cont-on-def continuous-on-eq-continuous-at
    flip: isCont-iff-restriction-cont-at)

```

### 3.2 Equivalence between Lipschitz Map and Restriction shift Map

For a function  $f :: 'a \Rightarrow 'b$ , it is equivalent to have *restriction-shift-on*  $f k A$  and *lipschitz-with-on*  $f (\delta_{\downarrow} TYPE('b))^k A$  when ' $a$ ' is of sort *geometric-restriction-space* and  $\sigma_{\downarrow} TYPE('b) = \sigma_{\downarrow} TYPE('a)$  ( $'b$  is then necessarily also of sort *geometric-restriction-space*).

Weaker versions of this result can be established with weaker

assumptions on the sort, this is what we do below.

```

lemma restriction-shift-nonneg-on-imp-lipschitz-with-on :
  ⟨lipschitz-with-on f (restriction-δ TYPE('b) ^ k) A⟩ if ⟨restriction-shift-on f (int k) A⟩
  and le-restriction-σ : ⟨ $\bigwedge n. \text{restriction-}\sigma \text{ TYPE('}b\text{'}) n \leq \text{restriction-}\sigma \text{ TYPE('}a\text{'}) n$ ⟩
for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
proof (rule lipschitz-with-onI)
  show ⟨ $0 \leq \text{restriction-}\delta \text{ TYPE('}b\text{'}) ^ k$ ⟩ by simp
next
  fix x y assume ⟨ $x \in A$ ⟩ ⟨ $y \in A$ ⟩ ⟨ $x \neq y$ ⟩ ⟨ $f x \neq f y$ ⟩
  from ⟨restriction-shift-on f k A⟩[THEN restriction-shift-onD, OF ⟨ $x \in A$ ⟩ ⟨ $y \in A$ ⟩]
  have ⟨ $i \in \text{restriction-related-set } x y \implies i + k \in \text{restriction-related-set } (f x) (f y)$ ⟩ for i
    by (simp add: nat-int-add)
  hence ⟨ $\text{Sup } (\text{restriction-related-set } x y) + k \in \text{restriction-related-set } (f x) (f y)$ ⟩
    using ⟨ $x \neq y$ ⟩ Sup-in-restriction-related-set by blast
  hence ⟨ $\text{Sup } (\text{restriction-related-set } x y) + k \leq \text{Sup } (\text{restriction-related-set } (f x) (f y))$ ⟩
    by (simp add: ⟨ $f x \neq f y$ ⟩ bdd-above-restriction-related-set-iff cSup-upper)
  moreover have ⟨ $\text{dist } (f x) (f y) = \text{restriction-}\sigma \text{ TYPE('}b\text{'}) (\text{Sup } (\text{restriction-related-set } (f x) (f y)))$ ⟩
    by (simp add: ⟨ $f x \neq f y$ ⟩ dist-restriction-is-bis-simplified)
  ultimately have ⟨ $\text{dist } (f x) (f y) \leq \text{restriction-}\sigma \text{ TYPE('}b\text{'}) (\text{Sup } (\text{restriction-related-set } x y) + k)$ ⟩
    by (simp add: decseqD decseq-restriction-space-class.decseq-restriction-σ)
  hence ⟨ $\text{dist } (f x) (f y) \leq \text{restriction-}\sigma \text{ TYPE('}b\text{'}) (\text{Sup } (\text{restriction-related-set } x y)) * \text{restriction-}\delta \text{ TYPE('}b\text{'}) ^ k$ ⟩
    by (meson order.trans restriction-σ-le-restriction-σ-times-pow-restriction-δ)
  also have ⟨ $\dots \leq \text{restriction-}\sigma \text{ TYPE('}a\text{'}) (\text{Sup } (\text{restriction-related-set } x y)) * \text{restriction-}\delta \text{ TYPE('}b\text{'}) ^ k$ ⟩
    by (simp add: le-restriction-σ)
  finally show ⟨ $\text{dist } (f x) (f y) \leq \text{restriction-}\delta \text{ TYPE('}b\text{'}) ^ k * \text{dist } x y$ ⟩
    by (simp add: dist-restriction-is-bis[of x y] ⟨ $x \neq y$ ⟩ mult.commute)
qed

corollary restriction-shift-nonneg-imp-lipschitz-with :
  ⟨[restriction-shift f (int k);  $\bigwedge n. \text{restriction-}\sigma \text{ TYPE('}b\text{'}) n \leq \text{restriction-}\sigma \text{ TYPE('}a\text{'}) n$ ] ⟩
   $\implies$  lipschitz-with f (restriction-δ TYPE('b) ^ k)
for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
using restriction-shift-def restriction-shift-nonneg-on-imp-lipschitz-with-on
by blast

```

```

lemma lipschitz-with-on-imp-restriction-shift-neg-on :
  ⟨restriction-shift-on f (− int k) A⟩ if ⟨lipschitz-with-on f (restriction-δ
  TYPE('b) powi − int k) A⟩
    and le-restriction-σ : ⟨ $\bigwedge n.$  restriction-σ TYPE('a)  $n \leq$  restriction-σ
  TYPE('b)  $n$ ⟩
for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
proof (rule restriction-shift-onI, goal-cases)
  fix x y n assume ⟨x ∈ A⟩ ⟨y ∈ A⟩ ⟨f x ≠ f y⟩ ⟨x ↓ n = y ↓ n⟩
  from ⟨f x ≠ f y⟩ have ⟨x ≠ y⟩ by blast
  from ⟨lipschitz-with-on f (restriction-δ TYPE('b) powi − int k) A⟩
    [THEN lipschitz-with-onD2, OF ⟨x ∈ A⟩ ⟨y ∈ A⟩]
  have ⟨dist (f x) (f y) ≤ restriction-δ TYPE('b) powi − int k * dist
  x y⟩ .
  hence ⟨dist (f x) (f y) * restriction-δ TYPE('b) ^ k ≤ dist x y⟩
    by (subst (asm) mult.commute)
    (drule mult-imp-div-pos-le[rotated]; simp add: power-int-minus-divide)
  hence ⟨restriction-σ TYPE('b) (Sup (restriction-related-set (f x) (f
  y))) * restriction-δ TYPE('b) ^ k
    ≤ restriction-σ TYPE('b) (Sup (restriction-related-set x y))⟩
    by (simp add: dist-restriction-is-bis ⟨x ≠ y⟩ ⟨f x ≠ f y⟩ )
    (drule order-trans[OF - le-restriction-σ], simp)
  from order-trans[OF restriction-σ-le-restriction-σ-times-pow-restriction-δ
    [of Sup (restriction-related-set (f x) (f y)) k this]]
  have ⟨restriction-σ TYPE('b) (Sup (restriction-related-set (f x) (f
  y))) + k
    ≤ restriction-σ TYPE('b) (Sup (restriction-related-set x y))⟩ .
  hence ⟨Sup (restriction-related-set x y) ≤ Sup (restriction-related-set
  (f x) (f y)) + k
    using strict-decseq-def-ter strict-decseq-restriction-σ by blast
    from order-trans[OF cSup-upper this] have ⟨n ≤ Sup (restriction-related-set
  (f x) (f y)) + k
    by (simp add: ⟨x ↓ n = y ↓ n⟩ ⟨x ≠ y⟩ bdd-above-restriction-related-set-iff)
    hence ⟨nat (int n + − int k) ≤ Sup (restriction-related-set (f x) (f
  y))⟩ by linarith
    thus ⟨f x ↓ nat (int n + − int k) = f y ↓ nat (int n + − int k)⟩
    by (metis not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified(2))
  qed

```

```

corollary lipschitz-with-imp-restriction-shift-neg :
  ⟨[lipschitz-with f (restriction-δ TYPE('b) powi − int k);
   $\bigwedge n.$  restriction-σ TYPE('a)  $n \leq$  restriction-σ TYPE('b)  $n]$ 
   $\implies$  restriction-shift f (− int k)⟩
for f :: ⟨'a :: decseq-restriction-space  $\Rightarrow$  'b :: at-least-geometric-restriction-space⟩
using lipschitz-with-on-imp-restriction-shift-neg-on restriction-shift-def
by blast

```

We obtained that *restriction-shift* implies *lipschitz-with* when  $0 \leq k$  and that *lipschitz-with* implies *restriction-shift* when  $k \leq 0$ . To cover the remaining cases, we give move from *at-least-geometric-restriction-space*

to geometric-restriction-space.

```

theorem lipschitz-with-on-iff-restriction-shift-on :
  ⟨lipschitz-with-on f (restriction-δ TYPE('b) powi k) A ⟷ restriction-shift-on f k A⟩
  if same-restriction-σ : ⟨restriction-σ TYPE('b) = restriction-σ TYPE('a)⟩
  for f :: ⟨'a :: decseq-restriction-space ⇒ 'b :: geometric-restriction-space⟩
proof (rule iff)
  — We could do a case on k, but both cases are actually handled by
  the proof required after applying [lipschitz-with-on ?f (δ↓ TYPE(?'b)
  powi - int ?k) ?A; ∏n. σ↓ TYPE(?'a) n ≤ σ↓ TYPE(?'b) n] ⟹
  restriction-shift-on ?f (- int ?k) ?A.
  show ⟨restriction-shift-on f k A⟩ if ⟨lipschitz-with-on f (restriction-δ
  TYPE('b) powi k) A⟩
  proof (rule restriction-shift-onI)
    fix x y n assume ⟨x ∈ A⟩ ⟨y ∈ A⟩ ⟨f x ≠ f y⟩ ⟨x ↓ n = y ↓ n⟩
    from ⟨f x ≠ f y⟩ have ⟨x ≠ y⟩ by blast
    from ⟨lipschitz-with-on f (restriction-δ TYPE('b) powi k) A⟩
      [THEN lipschitz-with-onD2, OF ⟨x ∈ A⟩ ⟨y ∈ A⟩]
    have ⟨dist (f x) (f y) ≤ restriction-δ TYPE('b) powi k * dist x y⟩ .
    also have ⟨dist (f x) (f y) = restriction-δ TYPE('b) ^ Sup
    (restriction-related-set (f x) (f y))⟩
      by (simp add: dist-restriction-is-bis ⟨f x ≠ f y⟩ restriction-σ-is)
    also have ⟨dist x y = restriction-δ TYPE('b) ^ Sup (restriction-related-set
    x y)⟩
      by (simp add: dist-restriction-is-bis ⟨x ≠ y⟩ restriction-σ-is
    same-restriction-σ[symmetric])
    finally have ⟨restriction-δ TYPE('b) ^ Sup (restriction-related-set
    (f x) (f y))⟩
      ≤ restriction-δ TYPE('b) powi k * restriction-δ TYPE('b)
      ^ Sup (restriction-related-set x y) .
    also have ⟨... = restriction-δ TYPE('b) powi (k + Sup (restriction-related-set
    x y))⟩
      by (subst power-int-add, metis order-less-irrefl zero-less-restriction-δ)
    simp
    finally have ⟨restriction-δ TYPE('b) ^ Sup (restriction-related-set
    (f x) (f y)) ≤ ...⟩ .
    moreover have ⟨restriction-δ TYPE('b) powi m ≤ restriction-δ
    TYPE('b) powi m' ⟹ m' ≤ m⟩ for m m'
      by (rule ccontr, unfold not-le,
        drule power-int-strict-decreasing[OF - zero-less-restriction-δ
        restriction-δ-less-one])
        (fold not-le, blast)
    ultimately have ⟨k + Sup (restriction-related-set x y) ≤ Sup
    (restriction-related-set (f x) (f y))⟩ by simp
    hence ⟨Sup (restriction-related-set x y) ≤ Sup (restriction-related-set
    (f x) (f y)) - k⟩ by simp
    with ⟨x ↓ n = y ↓ n⟩ ⟨x ≠ y⟩ linorder-not-le
    not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified(3)
    have ⟨n ≤ Sup (restriction-related-set (f x) (f y)) - k⟩ by fastforce
  
```

```

  hence ⟨nat (int n + k) ≤ Sup (restriction-related-set (f x) (f y))⟩
  by simp
    thus ⟨f x ↓ nat (int n + k) = f y ↓ nat (int n + k)⟩
      by (metis not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified(2))
    qed
  next
    show ⟨lipschitz-with-on f (restriction-δ TYPE('b) powi k) A⟩ if
      ⟨restriction-shift-on f k A⟩
      proof (cases k)
        show ⟨k = int k' ⟹ lipschitz-with-on f (restriction-δ TYPE('b)
          powi k) A⟩ for k'
          by (simp, rule restriction-shift-nonneg-on-imp-lipschitz-with-on)
            (use ⟨restriction-shift-on f k A⟩ same-restriction-σ in simp-all)
      next
        fix k' assume ⟨k = - int (Suc k')⟩
        show ⟨lipschitz-with-on f (restriction-δ TYPE('b) powi k) A⟩
        proof (rule lipschitz-with-onI)
          show ⟨0 ≤ restriction-δ TYPE('b) powi k⟩ by simp
        next
          fix x y assume ⟨x ∈ A⟩ ⟨y ∈ A⟩ ⟨x ≠ y⟩ ⟨f x ≠ f y⟩
          have ⟨i ∈ restriction-related-set x y ⟹ i = Suc k' ∈ restriction-related-set (f x) (f y)⟩ for i
            using ⟨restriction-shift-on f k A⟩[THEN restriction-shift-onD,
              OF ⟨x ∈ A⟩ ⟨y ∈ A⟩, of i]
            by (unfold ⟨k = - int (Suc k')⟩, clarify) (metis diff-conv-add-uminus
              nat-minus-as-int)
            hence ⟨Sup (restriction-related-set x y) = Suc k' ∈ restriction-related-set (f x) (f y)⟩
            using ⟨x ≠ y⟩ not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified
            by blast
            hence ⟨Sup (restriction-related-set x y) = Suc k' ≤ Sup (restriction-related-set
              (f x) (f y))⟩
              by (simp add: ⟨f x ≠ f y⟩ bdd-above-restriction-related-set-iff
                cSup-upper)
            hence * : ⟨Sup (restriction-related-set x y) + k ≤ Sup (restriction-related-set
              (f x) (f y))⟩
              by (simp add: ⟨k = - int (Suc k')⟩)
              have ⟨dist (f x) (f y) = restriction-δ TYPE('b) powi Sup
                (restriction-related-set (f x) (f y))⟩
                by (simp add: ⟨f x ≠ f y⟩ dist-restriction-is-bis-simplified
                  restriction-σ-is)
              also have ⟨... ≤ restriction-δ TYPE('b) powi (Sup (restriction-related-set
                x y) + k)⟩
                by (rule power-int-decreasing[OF *]; simp)
                  (metis order-less-irrefl zero-less-restriction-δ)
              also have ⟨... = restriction-δ TYPE('b) powi k * restriction-δ
                TYPE('b) powi (Sup (restriction-related-set x y))⟩
                by (subst power-int-add, metis order-less-irrefl zero-less-restriction-δ)
              simp

```

```

also have <restriction- $\delta$  TYPE('b) powi (Sup (restriction-related-set
x y)) = dist x y>
  by (simp add: <x ≠ y> dist-restriction-is-bis-simplified
    restriction- $\sigma$ -is same-restriction- $\sigma$ [symmetric])
  finally show <dist (f x) (f y) ≤ restriction- $\delta$  TYPE('b) powi k *
    dist x y .
```

qed  
qed  
qed

**corollary** lipschitz-with-iff-restriction-shift :

```

<restriction- $\sigma$  TYPE('b) = restriction- $\sigma$  TYPE('a) ==>
  lipschitz-with f (restriction- $\delta$  TYPE('b) powi k)  $\longleftrightarrow$  restriction-shift
  f k>
  for f :: <'a :: decseq-restriction-space  $\Rightarrow$  'b :: geometric-restriction-space>
  by (simp add: lipschitz-with-on-iff-restriction-shift-on restriction-shift-def)
```

## 4 Functions

### 4.1 Restriction Space

**instantiation** <fun> :: (type, restriction- $\sigma$ ) restriction- $\sigma$   
**begin**

**definition** restriction- $\sigma$ -fun :: <('a  $\Rightarrow$  'b) itself  $\Rightarrow$  nat  $\Rightarrow$  real>  
**where** <restriction- $\sigma$ -fun -  $\equiv$  restriction- $\sigma$  TYPE('b)>

**instance by** intro-classes

**end**

**instantiation** <fun> :: (type, non-decseq-restriction-space) non-decseq-restriction-space  
**begin**

**definition** dist-fun :: <['a  $\Rightarrow$  'b, 'a  $\Rightarrow$  'b]  $\Rightarrow$  real>  
**where** <dist-fun fg  $\equiv$  INF n  $\in$  restriction-related-set fg. restriction- $\sigma$   
 TYPE('a  $\Rightarrow$  'b) n>

**definition** uniformity-fun :: <((('a  $\Rightarrow$  'b)  $\times$  ('a  $\Rightarrow$  'b)) filter>  
**where** <uniformity-fun  $\equiv$  INF e  $\in$  {0 <..}. principal {(x, y). dist x y  
 < e}>

**definition** open-fun :: <('a  $\Rightarrow$  'b) set  $\Rightarrow$  bool>  
**where** <open-fun U  $\equiv$   $\forall$  x  $\in$  U. eventually ( $\lambda$ (x', y). x' = x  $\longrightarrow$  y  $\in$   
 U) uniformity>

```

instance by intro-classes
  (simp-all add: restriction- $\sigma$ -fun-def dist-fun-def open-fun-def
  uniformity-fun-def restriction- $\sigma$ -tendsto-zero)

end

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{decseq-restriction-space})$  decseq-restriction-space
  by intro-classes (simp add: restriction- $\sigma$ -fun-def decseq-restriction- $\sigma$ )

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{strict-decseq-restriction-space})$  strict-decseq-restriction-space
  by intro-classes
    (simp add: restriction- $\sigma$ -fun-def strict-decseq-restriction- $\sigma$ )

instantiation  $\langle \text{fun} \rangle :: (\text{type}, \text{restriction-}\delta)$  restriction- $\delta$ 
begin

definition restriction- $\delta$ -fun ::  $\langle ('a \Rightarrow 'b) \text{ itself} \Rightarrow \text{real} \rangle$ 
  where  $\langle \text{restriction-}\delta\text{-fun} - \equiv \text{restriction-}\delta \text{ TYPE}('b) \rangle$ 

instance by intro-classes (simp-all add: restriction- $\delta$ -fun-def)

end

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{at-least-geometric-restriction-space})$  at-least-geometric-restriction-space
proof intro-classes
  show  $\langle 0 < \text{restriction-}\sigma \text{ TYPE}('a \Rightarrow 'b) n \rangle$  for  $n$ 
    by (simp add: restriction- $\sigma$ -fun-def)
next
  show  $\langle \text{restriction-}\sigma \text{ TYPE}('a \Rightarrow 'b) (\text{Suc } n) \leq \text{restriction-}\delta \text{ TYPE}('a \Rightarrow 'b) * \text{restriction-}\sigma \text{ TYPE}('a \Rightarrow 'b) \rangle$ 
   $n \rangle$  for  $n$ 
    by (simp add: restriction- $\sigma$ -le restriction- $\sigma$ -fun-def restriction- $\delta$ -fun-def)
next
  show  $\langle \text{dist } f g = \text{Inf} (\text{restriction-}\sigma\text{-related-set } f g) \rangle$  for  $f g :: ('a \Rightarrow 'b)$ 
  by (simp add: dist-fun-def)
qed

instance  $\langle \text{fun} \rangle :: (\text{type}, \text{geometric-restriction-space})$  geometric-restriction-space
proof intro-classes
  show  $\langle \text{restriction-}\sigma \text{ TYPE}('a \Rightarrow 'b) n = \text{restriction-}\delta \text{ TYPE}('a \Rightarrow 'b) \wedge n \rangle$  for  $n$ 
    by (simp add: restriction- $\sigma$ -fun-def restriction- $\sigma$ -is restriction- $\delta$ -fun-def)
next
  show  $\langle \text{dist } f g = \text{Inf} (\text{restriction-}\sigma\text{-related-set } f g) \rangle$  for  $f g :: ('a \Rightarrow 'b)$ 

```

```

    by (simp add: dist-fun-def)
qed

lemma dist-image-le-dist-fun : <dist (f x) (g x) ≤ dist f g>
  for f g :: <'a ⇒ 'b :: non-decseq-restriction-space>
proof (unfold dist-restriction-is, rule cInf-superset-mono)
  show <restriction-σ-related-set f g ≠ {}> by simp
next
  show <bdd-below (restriction-σ-related-set (f x) (g x))>
    by (simp add: bounded-imp-bdd-below bounded-restriction-σ-related-set)
next
  show <restriction-σ-related-set f g ⊆ restriction-σ-related-set (f x) (g
x)>
    unfolding restriction-fun-def restriction-σ-fun-def
    by (simp add: image-def subset-iff) metis
qed

```

```

lemma Sup-dist-image-le-dist-fun : <(SUP x. dist (f x) (g x)) ≤ dist f
g>
  for f g :: <'a ⇒ 'b :: non-decseq-restriction-space>
  by (simp add: dist-image-le-dist-fun cSUP-least)
  — The other inequality will require the additional assumption
decseq.

```

```

context fixes f g :: <'a ⇒ 'b :: decseq-restriction-space> begin

lemma reached-dist-fun : <∃ x. dist f g = dist (f x) (g x)>
proof (cases <f = g>)
  show <f = g ⇒ ∃ x. dist f g = dist (f x) (g x)> by simp
next
  assume <f ≠ g>
  let ?n = <Sup (restriction-related-set f g)>
  from not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified(2,
3)[OF <f ≠ g>]
  obtain x where <∀ m ≤ ?n. f x ↓ m = g x ↓ m> <f x ↓ Suc ?n ≠ g x
↓ Suc ?n>
    unfolding restriction-fun-def by (meson lessI)
    hence <dist (f x) (g x) = restriction-σ TYPE('b) ?n>
      by (metis (no-types, lifting) le-neq-implies-less not-less-eq-eq
        not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified)
    with not-eq-imp-dist-restriction-is-restriction-σ-Sup-restriction-eq-simplified(1)
      [OF <f ≠ g>, unfolded restriction-σ-fun-def]
    have <dist f g = dist (f x) (g x)> by simp
    thus <∃ x. dist f g = dist (f x) (g x)> ..
qed

```

```

lemma dist-fun-eq-Sup-dist-image : <dist f g = (SUP x. dist (f x) (g x))>
proof (rule order-antisym)
  show <(SUP x. dist (f x) (g x)) ≤ dist f g> by (fact Sup-dist-image-le-dist-fun)
next
  from reached-dist-fun obtain x where <dist f g = dist (f x) (g x)>
 $\dots$ 
  thus <dist f g ≤ (SUP x. dist (f x) (g x))>
  proof (rule ord-eq-le-trans)
    show <dist (f x) (g x) ≤ (SUP x. dist (f x) (g x))>
    proof (rule cSup-upper)
      show <dist (f x) (g x) ∈ range (λx. dist (f x) (g x))> by simp
    next
      show <bdd-above (range (λx. dist (f x) (g x)))>
      by (rule bdd-aboveI[of- <dist f g>]) (auto intro: dist-image-le-dist-fun)
    qed
  qed
  qed

```

```

lemma fun-restriction-space-Sup-properties :
  <dist (f x) (g x) ≤ dist f g>
  <(λx. dist (f x) (g x) ≤ b) ⇒ dist f g ≤ b>
  by (use reached-dist-fun in <auto simp add: dist-image-le-dist-fun>)

```

**end**

## 4.2 Completeness

Actually we can obtain even better: when the instance '*b*' of *decseq-restriction-space* is also an instance of *complete-space*, the type '*a* → '*b*' is an instance of *complete-space*.

This is because when '*b*' is an instance of *decseq-restriction-space* (and not only *non-decseq-restriction-space*) the distance between two functions is reached (see  $\exists x. \text{dist } ?f ?g = \text{dist} (?f x) (?g x)$ ).

The only remaining thing is to prove that completeness is preserved on higher-order.

```

instance <fun> :: (type, complete-decseq-restriction-space) complete-decseq-restriction-space
  by intro-classes (fact restriction-chain-imp-restriction-convergent)

instance <fun> :: (type, complete-strict-decseq-restriction-space) complete-strict-decseq-restriction-space
  by intro-classes (fact restriction-chain-imp-restriction-convergent)

```

```

instance <fun> :: (type, complete-at-least-geometric-restriction-space)
complete-at-least-geometric-restriction-space
by intro-classes (fact restriction-chain-imp-restriction-convergent)

instance <fun> :: (type, complete-geometric-restriction-space) complete-geometric-restriction-space
by intro-classes (fact restriction-chain-imp-restriction-convergent)

```

### 4.3 Kind of Extensionality

```

context fixes f :: <['a :: metric-space, 'b :: type] ⇒
'c :: deseq-restriction-space> begin

lemma lipschitz-with-simplification:
<lipschitz-with f α ↔ (forall y. lipschitz-with (λx. f x y) α)>
proof (intro iffI allI)
fix y assume assm : <lipschitz-with f α>
show <lipschitz-with (λx. f x y) α>
proof (rule lipschitz-withI)
from assm[THEN lipschitz-withD1] show <0 ≤ α> .
next
show <dist (f x1 y) (f x2 y) ≤ α * dist x1 x2> for x1 x2
by (rule order-trans[OF - assm[THEN lipschitz-withD2]])
(simp add: dist-image-le-dist-fun)
qed
next
assume assm : <forall y. lipschitz-with (λx. f x y) α>
show <lipschitz-with f α>
proof (rule lipschitz-withI)
from assm[rule-format, THEN lipschitz-withD1] show <0 ≤ α> .
next
fix x1 x2
obtain y where <dist (f x1) (f x2) = dist (f x1 y) (f x2 y)>
by (meson reached-dist-fun)
also have <... ≤ α * dist x1 x2> by (rule assm[rule-format, THEN
lipschitz-withD2])
finally show <dist (f x1) (f x2) ≤ α * dist x1 x2> .
qed
qed

```

```

lemma non-expanding-simplification :
<non-expanding f ↔ (forall y. non-expanding (λx. f x y))>
by (metis lipschitz-with-simplification non-expanding-on-def)

```

```

lemma contraction-with-simplification:
<contraction-with f α ↔ (forall y. contraction-with (λx. f x y) α)>
by (metis contraction-with-on-def lipschitz-with-simplification)

```

```
end
```

## 5 Product

The product type ' $a \times b$ ' of two metric spaces is already instantiated as a metric space by setting  $dist x y = sqrt ((dist (fst x) (fst y))^2 + (dist (snd x) (snd y))^2)$ . Unfortunately, this definition is not compatible with the distance required by the *non-decseq-restriction-space*. We first have to define a new product type with a trivial **typedef**.

### 5.1 Isomorphic Product Construction

#### 5.1.1 Definition and First Properties

```
typedef ('a, 'b) prodmax ((- ×max / -) [21, 20] 20) = <UNIV :: ('a × 'b) set>
morphisms from-prodmax to-prodmax by simp
```

— Simplifications because the **typedef** is trivial.

```
declare from-prodmax-inject [simp]
from-prodmax-inverse [simp]
```

```
lemmas to-prodmax-inject-simplified [simp] = to-prodmax-inject [simplified]
and to-prodmax-inverse-simplified [simp] = to-prodmax-inverse [simplified]
```

```
lemmas to-prodmax-induct-simplified = to-prodmax-induct [simplified]
and to-prodmax-cases-simplified = to-prodmax-cases [simplified]
and from-prodmax-induct-simplified = from-prodmax-induct [simplified]
and from-prodmax-cases-simplified = from-prodmax-cases [simplified]
```

```
setup-lifting type-definition-prodmax
```

```
lift-definition Pairmax :: <'a ⇒ 'b ⇒ 'a ×max 'b> is Pair .
```

```
free-constructors case-prodmax for Pairmax fstmax sndmax
by (metis Pairmax.abs-eq from-prodmax-inverse surjective-pairing)
  (metis Pairmax.rep-eq prod.inject)
```

```
lemma fstmax-def : fstmax ≡ map-fun from-prodmax id fst
by (intro eq-reflection ext, simp add: fstmax-def,
```

```

metis Pairmax.rep-eq from-prodmax-inverse fstmax-def prod.collapse
prodmax.sel(1)
lemma fstmax-rep-eq : <fstmax x = fst (from-prodmax x)>
by (metis Pairmax.rep-eq fst-conv prodmax.collapse)

lemma fstmax-abs-eq [simp] : <fstmax (to-prodmax y) = fst y>
by (metis Pairmax.abs-eq prod.exhaust-sel prodmax.sel(1))

lemma fstmax-transfer [transfer-rule] : <rel-fun (pcr-prodmax (=) (=))
(=) fst fstmax>
by (metis (mono-tags) Pairmax.rep-eq cr-prodmax-def fst-conv prodmax.collapse
prodmax.pcr-cr-eq rel-funI)

lemma sndmax-def : <sndmax ≡ map-fun from-prodmax id snd>
by (intro eq-reflection ext, simp add: sndmax-def,
      metis Pairmax.rep-eq from-prodmax-inverse prod.collapse prodmax.case)

lemma sndmax-rep-eq : <sndmax x = snd (from-prodmax x)>
by (metis Pairmax.rep-eq prodmax.collapse snd-conv)

lemma sndmax-abs-eq [simp] : <sndmax (to-prodmax y) = snd y>
by (metis Pairmax.abs-eq prod.exhaust-sel prodmax.sel(2))

lemma sndmax-transfer [transfer-rule] : <rel-fun (pcr-prodmax (=)
(=)) (=) snd sndmax>
by (metis (mono-tags, lifting) Pairmax.rep-eq cr-prodmax-def
prodmax.collapse prodmax.pcr-cr-eq rel-funI snd-conv)

lemma case-prodmax-def : <case-prodmax ≡ map-fun id (map-fun
from-prodmax id) case-prod>
by (intro eq-reflection ext, simp add: prodmax.case-eq-if fstmax-rep-eq
sndmax-rep-eq split-beta)

lemma case-prodmax-rep-eq : <case-prodmax f p = (case from-prodmax
p of (x, y) ⇒ f x y)>
by (simp add: fstmax-rep-eq prodmax.case-eq-if sndmax-rep-eq split-beta)

lemma case-prodmax-abs-eq [simp] : <case-prodmax f (to-prodmax q)
= (case q of (x, y) ⇒ f x y)>
by (simp add: prodmax.case-eq-if split-beta)

lemma case-prodmax-transfer [transfer-rule] : <rel-fun (=) (rel-fun
(pcr-prodmax (=) (=)) (=)) case-prod case-prodmax>
by (simp add: cr-prodmax-def fstmax-rep-eq prodmax.case-eq-if prodmax.pcr-cr-eq
rel-fun-def sndmax-rep-eq split-beta)

```

## 5.2 Syntactic Sugar

The following syntactic sugar is of course recovered from the theory *HOL.Product-Type*.

```
nonterminal tuple-argsmax and patternsmax
syntax
-tuplemax :: 'a ⇒ tuple-argsmax ⇒ 'a ×max 'b      ((1⟨-, / -⟩)⟨)
-tuple-argmax :: 'a ⇒ tuple-argsmax                  (↔)
-tuple-argsmax :: 'a ⇒ tuple-argsmax ⇒ tuple-argsmax  (⟨-, / -⟩)
-patternmax   :: pttrn ⇒ patternsmax ⇒ pttrn       ((⟨-, / -⟩)⟨)
              :: pttrn ⇒ patternsmax                   (↔)
-patternsmax :: pttrn ⇒ patternsmax ⇒ patternsmax  (⟨-, / -⟩)
translations
⟨x, y⟩ ⇌ CONST Pairmax x y
-patternmax x y ⇌ CONST Pairmax x y
-patternsmax x y ⇌ CONST Pairmax x y
-tuplemax x (-tuple-argsmax y z) ⇌ -tuplemax x (-tuple-argmax
(-tuplemax y z))
λ⟨x, y, zs⟩. b ⇌ CONST case-prodmax (λx ⟨y, zs⟩. b)
λ⟨x, y⟩. b ⇌ CONST case-prodmax (λx y. b)
-abs (CONST Pairmax x y) t → λ⟨x, y⟩. t
— This rule accommodates tuples in case C ... ⟨x, y⟩ ... ⇒ ...:
The ⟨x, y⟩ is parsed as Pairmax x y because it is logic, not pttrn.
```

With this syntactic sugar, one can write *case a of* ⟨b, c, d, e⟩ ⇒ ⟨c, d⟩, λ⟨y, u⟩. a, λ⟨a, b⟩. ⟨a, b, c, d, e⟩, λ⟨a, b, c⟩. a, ... as for the type 'a × 'b'.

```
lemmas to-prodmax-tuple [simp] = Pairmax.abs-eq[symmetric]
and from-prodmax-tuplemax [simp] = Pairmax.rep-eq
```

## 5.3 Product

We first redo the work of *Restriction-Spaces.Restriction-Spaces-Prod*.

```
instantiation prodmax :: (restriction, restriction) restriction
begin
```

```
lift-definition restriction-prodmax :: ⟨'a ×max 'b ⇒ nat ⇒ 'a ×max
'b⟩ is ⟨(↓)⟩ .
```

```
lemma restriction-prodmax-def' : ⟨p ↓ n = ⟨fstmax p ↓ n, sndmax p
↓ n⟩⟩
by transfer (simp add: restriction-prod-def)
```

```
instance by (intro-classes, transfer, simp)
```

```
end
```

```

instance prodmax :: (restriction-space, restriction-space) restriction-space
  by (intro-classes; transfer) (simp-all add: ex-not-restriction-related)

```

```

instantiation prodmax :: (restriction- $\sigma$ , restriction- $\sigma$ ) restriction- $\sigma$ 
begin

```

```

definition restriction- $\sigma$ -prodmax :: <('a  $\times_{max}$  'b) itself  $\Rightarrow$  nat  $\Rightarrow$  real>
  where <restriction- $\sigma$ -prodmax - n  $\equiv$ 
    max (restriction- $\sigma$  TYPE('a) n) (restriction- $\sigma$  TYPE('b) n)>

```

```

instance by intro-classes
end

```

```

instantiation prodmax :: (non-decseq-restriction-space, non-decseq-restriction-space)
  non-decseq-restriction-space
begin

```

```

definition dist-prodmax :: <['a  $\times_{max}$  'b, 'a  $\times_{max}$  'b]  $\Rightarrow$  real>
  where <dist-prodmax f g  $\equiv$  INF n  $\in$  restriction-related-set f g. restriction- $\sigma$  TYPE('a  $\times_{max}$  'b) n>

```

```

definition uniformity-prodmax :: <((('a  $\times_{max}$  'b)  $\times$  'a  $\times_{max}$  'b) filter>
  where <uniformity-prodmax  $\equiv$  INF e $\in\{0 <..\}$ . principal {(x, y). dist x y < e}>

```

```

definition open-prodmax :: <('a  $\times_{max}$  'b) set  $\Rightarrow$  bool>
  where <open-prodmax U  $\equiv$   $\forall x \in U$ . eventually ( $\lambda(x', y)$ .  $x' = x \longrightarrow y \in U$ ) uniformity>

```

```

instance
proof intro-classes
  show <restriction- $\sigma$  TYPE('a  $\times_{max}$  'b)  $\longrightarrow$  0>
    by (rule real-tendsto-sandwich
      [of < $\lambda n$ . 0> - - < $\lambda n$ . restriction- $\sigma$  TYPE('a) n + restriction- $\sigma$  TYPE('b) n>])
      (simp-all add: order-less-imp-le restriction- $\sigma$ -prodmax-def max-def
        restriction- $\sigma$ -tendsto-zero tendsto-add-zero)
  qed (simp-all add: uniformity-prodmax-def open-prodmax-def
    restriction- $\sigma$ -prodmax-def max-def dist-prodmax-def)

```

```

end

```

```

instance prodmax :: (decseq-restriction-space, decseq-restriction-space)
  decseq-restriction-space
proof intro-classes

```

```

show ⟨decseq (restriction- $\sigma$  TYPE('a  $\times_{max}$  'b))⟩
proof (intro decseq-SucI)
  show ⟨restriction- $\sigma$  TYPE('a  $\times_{max}$  'b) (Suc n)  $\leq$  restriction- $\sigma$ 
  TYPE('a  $\times_{max}$  'b) n⟩ for n
    using decseq-SucD[of ⟨restriction- $\sigma$  TYPE('a)⟩ n]
    decseq-SucD[of ⟨restriction- $\sigma$  TYPE('b)⟩ n]
    by (auto simp add: restriction- $\sigma$ -prodmax-def decseq-restriction- $\sigma$ )
  qed
qed

instance prodmax :: (strict-decseq-restriction-space, strict-decseq-restriction-space)
strict-decseq-restriction-space
proof intro-classes
  show ⟨strict-decseq (restriction- $\sigma$  TYPE('a  $\times_{max}$  'b))⟩
  proof (intro strict-decseq-SucI)
    show ⟨restriction- $\sigma$  TYPE('a  $\times_{max}$  'b) (Suc n) < restriction- $\sigma$ 
    TYPE('a  $\times_{max}$  'b) n⟩ for n
      using strict-decseq-SucD[of ⟨restriction- $\sigma$  TYPE('a)⟩ n]
      strict-decseq-SucD[of ⟨restriction- $\sigma$  TYPE('b)⟩ n]
      by (auto simp add: restriction- $\sigma$ -prodmax-def strict-decseq-restriction- $\sigma$ )
    qed
  qed

instantiation prodmax :: (restriction- $\delta$ , restriction- $\delta$ ) restriction- $\delta$ 
begin

definition restriction- $\delta$ -prodmax :: ⟨('a  $\times_{max}$  'b) itself  $\Rightarrow$  real⟩
  where ⟨restriction- $\delta$ -prodmax -  $\equiv$  max (restriction- $\delta$  TYPE('a))
  (restriction- $\delta$  TYPE('b))⟩

instance by intro-classes (simp-all add: restriction- $\delta$ -prodmax-def max-def)

end

instance prodmax :: (at-least-geometric-restriction-space, at-least-geometric-restriction-space)
at-least-geometric-restriction-space
proof intro-classes
  show ⟨0 < restriction- $\sigma$  TYPE('a  $\times_{max}$  'b) n⟩ for n by simp
next
  show ⟨restriction- $\sigma$  TYPE('a  $\times_{max}$  'b) (Suc n)
   $\leq$  restriction- $\delta$  TYPE('a  $\times_{max}$  'b) * restriction- $\sigma$  TYPE('a
   $\times_{max}$  'b) n⟩ for n
    by (auto intro: order-trans[OF restriction- $\sigma$ -le]
    simp add: restriction- $\delta$ -prodmax-def mult-mono' restriction- $\sigma$ -prodmax-def)
next
  show ⟨dist p1 p2 = Inf (restriction- $\sigma$ -related-set p1 p2)⟩ for p1 p2
  :: ⟨'a  $\times_{max}$  'b⟩
    by (simp add: dist-prodmax-def)

```

**qed**

```

instance prodmax :: (geometric-restriction-space, geometric-restriction-space)
geometric-restriction-space
proof intro-classes
  show ⟨restriction- $\sigma$  TYPE('a  $\times_{max}$  'b) n = restriction- $\delta$  TYPE('a
 $\times_{max}$  'b)  $\wedge$  n⟩ for n
    by (simp add: restriction- $\sigma$ -prodmax-def restriction- $\sigma$ -is restriction- $\delta$ -prodmax-def max-def)
      (meson nle-le power-mono zero-le-restriction- $\delta$ )
next
  show ⟨dist p1 p2 = Inf (restriction- $\sigma$ -related-set p1 p2)⟩ for p1 p2
  :: ⟨'a  $\times_{max}$  'b⟩
    by (simp add: dist-prodmax-def)
qed

```

```

lemma max-dist-projections-le-dist-prodmax :
  ⟨max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1) (sndmax p2)))
  ≤ dist p1 p2⟩
proof (unfold dist-restriction-is max-def, split if-split, intro conjI impI)
  show ⟨Inf (restriction- $\sigma$ -related-set (sndmax p1) (sndmax p2)) ≤ Inf
  (restriction- $\sigma$ -related-set p1 p2)⟩
    proof (rule cINF-superset-mono[OF nonempty-restriction-related-set])
      show ⟨bdd-below (restriction- $\sigma$ -related-set (sndmax p1) (sndmax
      p2))⟩
        by (meson bdd-belowI2 zero-le-restriction- $\sigma$ )
      qed (simp-all add: subset-iff add: restriction-prodmax-def' restriction- $\sigma$ -prodmax-def)
next
  show ⟨Inf (restriction- $\sigma$ -related-set (fstmax p1) (fstmax p2)) ≤ Inf
  (restriction- $\sigma$ -related-set p1 p2)⟩
    proof (rule cINF-superset-mono[OF nonempty-restriction-related-set])
      show ⟨bdd-below (restriction- $\sigma$ -related-set (fstmax p1) (fstmax p2))⟩
        by (meson bdd-belowI2 zero-le-restriction- $\sigma$ )
      qed (simp-all add: subset-iff add: restriction-prodmax-def' restriction- $\sigma$ -prodmax-def)
qed

```

## 5.4 Completeness

### 5.4.1 Preliminaries

**default-sort** non-decseq-restriction-space — Otherwise we should always specify.

```

lemma restriction- $\sigma$ -prodmax-commute :
  ⟨restriction- $\sigma$  TYPE('b  $\times_{max}$  'a) = restriction- $\sigma$  TYPE('a  $\times_{max}$ 

```

```

'b)›
  unfolding restriction-σ-prodmax-def by (rule ext) simp

definition dist-left-prodmax :: ⟨('a ×max 'b) itself ⇒ 'a ⇒ 'a ⇒ real⟩
  where ⟨dist-left-prodmax - x y ≡ INF n ∈ restriction-related-set x y. restriction-σ TYPE('a ×max 'b) n⟩

definition dist-right-prodmax :: ⟨('a ×max 'b) itself ⇒ 'b ⇒ 'b ⇒ real⟩
  where ⟨dist-right-prodmax - x y ≡ INF n ∈ restriction-related-set x y. restriction-σ TYPE('a ×max 'b) n⟩

lemma dist-right-prodmax-is-dist-left-prodmax :
  ⟨dist-right-prodmax TYPE('b ×max 'a) = dist-left-prodmax TYPE('a ×max 'b)⟩
  unfolding dist-left-prodmax-def dist-right-prodmax-def
  by (subst restriction-σ-prodmax-commute) simp

lemma dist-le-dist-left-prodmax : ⟨dist x y ≤ dist-left-prodmax TYPE('a ×max 'b) x y⟩
proof (unfold dist-left-prodmax-def dist-restriction-is,
       rule cINF-mono[OF nonempty-restriction-related-set[of x y]])
show ⟨bdd-below (restriction-σ-related-set x y)⟩
by (meson bdd-belowI2 zero-le-restriction-σ)
next
show ⟨m ∈ restriction-related-set x y ⟹
      ∃n∈restriction-related-set x y. σ↓ TYPE('a) n ≤ σ↓ TYPE('a ×max 'b) m⟩ for m
  by (metis max.cobounded1 restriction-σ-prodmax-def)
qed

lemma dist-le-dist-right-prodmax : ⟨dist x y ≤ dist-right-prodmax TYPE('b ×max 'a) x y⟩
by (simp add: dist-le-dist-left-prodmax dist-right-prodmax-is-dist-left-prodmax)

```

```

lemma
fixes p1 p2 :: ⟨'a :: decseq-restriction-space ×max 'b :: decseq-restriction-space⟩
shows dist-prodmax-le-max-dist-left-prodmax-dist-right-prodmax :
  ⟨dist p1 p2 ≤ max (dist-left-prodmax TYPE('a ×max 'b) (fstmax p1) (fstmax p2))
    (dist-right-prodmax TYPE('a ×max 'b) (sndmax p1) (sndmax p2))⟩
proof -
interpret left : DecseqRestrictionSpace ⟨(↓)⟩ ⟨(=)⟩ ⟨UNIV⟩
  ⟨restriction-σ TYPE('a ×max 'b)⟩ ⟨dist-left-prodmax TYPE('a ×max 'b)⟩

```

```

by unfold-locales
  (simp-all add: restriction-σ-tendsto-zero dist-left-prodmax-def
decseq-restriction-σ)

interpret right : DecseqRestrictionSpace ⟨(↓)⟩ ⟨(=)⟩ ⟨UNIV :: 'b set⟩
  ⟨restriction-σ TYPE('a ×max 'b)⟩ ⟨dist-right-prodmax TYPE('a
×max 'b)⟩
by unfold-locales
  (simp-all add: restriction-σ-tendsto-zero dist-right-prodmax-def
decseq-restriction-σ)

show ⟨dist p1 p2 ≤ max (dist-left-prodmax TYPE('a ×max 'b)
(fstmax p1) (fstmax p2))⟩
  ⟨dist-right-prodmax TYPE('a ×max 'b) (sndmax
p1) (sndmax p2))⟩
  by (auto simp add: dist-restriction-is-bis left.dist-restriction-is-bis
right.dist-restriction-is-bis prodmax.expand restriction-prodmax-def')
    (smt (verit, best) Collect-cong nle-le restriction-related-le)
qed

```

**default-sort** *type* — Back to normal.

#### 5.4.2 Complete Restriction Space

When the instances '*a*' and '*b*' of *decseq-restriction-space* are also instances of *complete-space*, the type '*a* ×<sub>max</sub> '*b*' is an instance of *complete-space*.

```

lemma restriction-chain-prodmax-iff :
  ⟨restriction-chain σ ⟷ restriction-chain (λn. fstmax (σ n)) ∧
    restriction-chain (λn. sndmax (σ n))⟩
  by (simp add: restriction-chain-def, transfer)
    (metis fst-conv prod.collapse restriction-prod-def snd-conv)

lemma restriction-tendsto-prodmax-iff :
  ⟨σ ↓→ Σ ⟷ (λn. fstmax (σ n)) ↓→ fstmax Σ ∧ (λn. sndmax
(σ n)) ↓→ sndmax Σ⟩
  by (simp add: restriction-tendsto-def, transfer, simp add: restriction-prod-def)
    (meson nle-le order.trans)

lemma restriction-convergent-prodmax-iff :
  ⟨restriction-convergent σ ⟷ restriction-convergent (λn. fstmax (σ
n)) ∧
    restriction-convergent (λn. sndmax (σ n))⟩
  by (simp add: restriction-convergent-def restriction-tendsto-prodmax-iff)
    (metis prodmax.sel)

```

```

instance prodmax :: (complete-decseq-restriction-space, complete-decseq-restriction-space)
  complete-decseq-restriction-space
proof (intro-classes, transfer)
  fix  $\sigma :: \langle \text{nat} \Rightarrow 'a \times_{\max} 'b \rangle$  assume  $\langle \text{chain}_{\downarrow} \sigma \rangle$ 
  hence  $\langle \text{chain}_{\downarrow} (\lambda n. \text{fst}_{\max} (\sigma n)) \rangle \langle \text{chain}_{\downarrow} (\lambda n. \text{snd}_{\max} (\sigma n)) \rangle$ 
    by (simp-all add: restriction-chain-prodmax-iff)
  hence  $\langle \text{convergent}_{\downarrow} (\lambda n. \text{fst}_{\max} (\sigma n)) \rangle \langle \text{convergent}_{\downarrow} (\lambda n. \text{snd}_{\max} (\sigma n)) \rangle$ 
    by (simp-all add: restriction-chain-imp-restriction-convergent)
  thus  $\langle \text{convergent}_{\downarrow} \sigma \rangle$ 
    by (simp add: restriction-convergent-prodmax-iff)
qed

```

```

instance prodmax :: (complete-strict-decseq-restriction-space, complete-strict-decseq-restriction-space)
  complete-strict-decseq-restriction-space
by intro-classes (simp add: restriction-chain-imp-restriction-convergent)

instance prodmax :: (complete-at-least-geometric-restriction-space, com-
  plete-at-least-geometric-restriction-space)
  complete-at-least-geometric-restriction-space
by intro-classes (simp add: restriction-chain-imp-restriction-convergent)

instance prodmax :: (complete-geometric-restriction-space, complete-geometric-restriction-space)
  complete-geometric-restriction-space
by intro-classes (simp add: restriction-chain-imp-restriction-convergent)

```

When the types ' $'a$ ' and ' $'b$ ' share the same restriction sequence, we have the following equality.

```

lemma same-restriction- $\sigma$ -imp-restriction- $\sigma$ -prodmax-is [simp] :
   $\langle \text{restriction-}\sigma \text{ TYPE}'('b :: \text{non-decseq-restriction-space}) =$ 
   $\text{restriction-}\sigma \text{ TYPE}'('a :: \text{non-decseq-restriction-space}) \Rightarrow$ 
   $\text{restriction-}\sigma \text{ TYPE}'('a \times_{\max} 'b) = \text{restriction-}\sigma \text{ TYPE}'('a)$ 
unfolding restriction- $\sigma$ -prodmax-def by simp

```

```

lemma same-restriction- $\sigma$ -imp-dist-prodmax-eq-max-dist-projections :
   $\langle \text{dist } p1 p2 = \max (\text{dist } (\text{fst}_{\max} p1) (\text{fst}_{\max} p2)) (\text{dist } (\text{snd}_{\max} p1) (\text{snd}_{\max} p2)) \rangle$ 
  if same-restriction- $\sigma$  [simp] :  $\langle \text{restriction-}\sigma \text{ TYPE}'('b) = \text{restriction-}\sigma \text{ TYPE}'('a) \rangle$ 
  for  $p1 p2 :: \langle 'a :: \text{decseq-restriction-space} \times_{\max} 'b :: \text{decseq-restriction-space} \rangle$ 
proof (rule order-antisym)
  have  $\langle \text{dist-left-prod}_{\max} \text{ TYPE}'('a \times_{\max} 'b) (\text{fst}_{\max} p1) (\text{fst}_{\max} p2) \rangle$ 
   $= \text{dist } (\text{fst}_{\max} p1) (\text{fst}_{\max} p2)$ 
  by (simp add: dist-left-prodmax-def dist-restriction-is)
  moreover have  $\langle \text{dist-right-prod}_{\max} \text{ TYPE}'('a \times_{\max} 'b) (\text{snd}_{\max} p1) (\text{snd}_{\max} p2) = \text{dist } (\text{snd}_{\max} p1) (\text{snd}_{\max} p2) \rangle$ 

```

```

by (simp add: dist-right-prodmax-def dist-restriction-is)
ultimately show ‹dist p1 p2 ≤ max (dist (fstmax p1) (fstmax p2))
(dist (sndmax p1) (sndmax p2))›
by (metis dist-prodmax-le-max-dist-left-prodmax-dist-right-prodmax)
next
show ‹max (dist (fstmax p1) (fstmax p2)) (dist (sndmax p1) (sndmax
p2)) ≤ dist p1 p2›
by (fact max-dist-projections-le-dist-prodmax)
qed

```

Now, one can write things like  $v \langle x, y \rangle . f \langle x, y \rangle$ .

We could of course imagine more support for ' $a \times_{max} b$ ' type, but as restriction spaces are intended to be used without recourse to metric spaces, we have not undertaken this task for the time being.

## 6 Main entry Point