

Relative Security

Andrei Popescu Jamie Wright

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Abstract

This entry formalizes the notion of relative security, which can be used to model transient execution vulnerabilities in the style of Spectre and Meltdown. The notion was introduced in the CSF 2023 paper “Relative Security: Formally Modeling and (Dis)Proving Resilience Against Semantic Optimization Vulnerabilities” by Brijesh Dongol, Matt Griffin, Andrei Popescu and Jamie Wright [1].

It defines two versions of relative security: a finitary one (restricted to finite traces), and an infinitary one (working with both finite and infinite traces). It formalizes unwinding methods for verifying relative security in both the finitary and infinitary versions, and proves their soundness. The proof of soundness in the infinitary case is a substantial application of Isabelle’s corecursion and coinduction infrastructure.

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1 Finitary Relative Security

This theory formalizes the finitary version of relative security, more precisely the notion expressed in terms of finite traces.

```
theory Relative-Security-fin
imports Preliminaries/Transition-System
begin
```

```
declare Let-def[simp]
```

```
no-notation relcomp (infixr O 75)
no-notation relcompp (infixr OO 75)
```

1.1 Finite-trace versions of leakage models and attacker models

```
locale Leakage-Mod-fin = System-Mod istate validTrans final
for istate :: 'state  $\Rightarrow$  bool and validTrans :: 'state  $\times$  'state  $\Rightarrow$  bool and final ::
'state  $\Rightarrow$  bool
+
fixes leakVia :: 'state list  $\Rightarrow$  'state list  $\Rightarrow$  'leak  $\Rightarrow$  bool
```

```
locale Attacker-Mod-fin = System-Mod istate validTrans final
for istate :: 'state  $\Rightarrow$  bool and validTrans :: 'state  $\times$  'state  $\Rightarrow$  bool and final ::
'state  $\Rightarrow$  bool
+
fixes S :: 'state list  $\Rightarrow$  'secret list
and A :: 'state trace  $\Rightarrow$  'act list
and O :: 'state trace  $\Rightarrow$  'obs list
begin
```

```
fun leakVia :: 'state list  $\Rightarrow$  'state list  $\Rightarrow$  'secret list  $\times$  'secret list  $\Rightarrow$  bool
where
leakVia tr tr' (sl,sl') = (S tr = sl  $\wedge$  S tr' = sl'  $\wedge$  A tr = A tr'  $\wedge$  O tr  $\neq$  O tr')
```

```
lemmas leakVia-def = leakVia.simps
```

```
end
```

```
sublocale Attacker-Mod-fin < Leakage-Mod-fin
where leakVia = leakVia
by standard
```

1.2 Locales for increasingly concrete notions of finitary relative security

locale *Relative-Security''-fin* =

Van: *Leakage-Mod-fin* *istateV* *validTransV* *finalV* *leakViaV*

+

Opt: *Leakage-Mod-fin* *istateO* *validTransO* *finalO* *leakViaO*
for *validTransV* :: 'stateV × 'stateV ⇒ bool

and *istateV* :: 'stateV ⇒ bool **and** *finalV* :: 'stateV ⇒ bool

and *leakViaV* :: 'stateV list ⇒ 'stateV list ⇒ 'leak ⇒ bool

and *validTransO* :: 'stateO × 'stateO ⇒ bool

and *istateO* :: 'stateO ⇒ bool **and** *finalO* :: 'stateO ⇒ bool

and *leakViaO* :: 'stateO list ⇒ 'stateO list ⇒ 'leak ⇒ bool

and *corrState* :: 'stateV ⇒ 'stateO ⇒ bool

begin

definition *rsecure* :: bool **where**

rsecure ≡ ∀ l s1 tr1 s2 tr2.

istateO s1 ∧ *Opt.validFromS* s1 tr1 ∧ *Opt.completedFrom* s1 tr1 ∧
istateO s2 ∧ *Opt.validFromS* s2 tr2 ∧ *Opt.completedFrom* s2 tr2 ∧
leakViaO tr1 tr2 l

→

(∃ sv1 trv1 sv2 trv2.

istateV sv1 ∧ *istateV* sv2 ∧ *corrState* sv1 s1 ∧ *corrState* sv2 s2 ∧
Van.validFromS sv1 trv1 ∧ *Van.completedFrom* sv1 trv1 ∧
Van.validFromS sv2 trv2 ∧ *Van.completedFrom* sv2 trv2 ∧
leakViaV trv1 trv2 l)

end

locale *Relative-Security'-fin* =

Van: *Attacker-Mod-fin* *istateV* *validTransV* *finalV* *SV* *AV* *OV*

+

Opt: *Attacker-Mod-fin* *istateO* *validTransO* *finalO* *SO* *AO* *OO*

for *validTransV* :: 'stateV × 'stateV ⇒ bool

and *istateV* :: 'stateV ⇒ bool **and** *finalV* :: 'stateV ⇒ bool

and *SV* :: 'stateV list ⇒ 'secret list

and *AV* :: 'stateV trace ⇒ 'actV list

and *OV* :: 'stateV trace ⇒ 'obsV list

and *validTransO* :: 'stateO × 'stateO ⇒ bool

and *istateO* :: 'stateO ⇒ bool **and** *finalO* :: 'stateO ⇒ bool

and *SO* :: 'stateO list ⇒ 'secret list

and *AO* :: 'stateO trace ⇒ 'actO list

and *OO* :: 'stateO trace ⇒ 'obsO list

and *corrState* :: 'stateV ⇒ 'stateO ⇒ bool

sublocale *Relative-Security'-fin* < *Relative-Security''-fin*
where *leakViaV* = *Van.leakVia* **and** *leakViaO* = *Opt.leakVia*
by *standard*

context *Relative-Security'-fin*
begin

lemma *rsecure-def2*:

rsecure \longleftrightarrow

($\forall s1\ tr1\ s2\ tr2.$

istateO *s1* \wedge *Opt.validFromS* *s1* *tr1* \wedge *Opt.completedFrom* *s1* *tr1* \wedge
istateO *s2* \wedge *Opt.validFromS* *s2* *tr2* \wedge *Opt.completedFrom* *s2* *tr2* \wedge
AO *tr1* = *AO* *tr2* \wedge *OO* *tr1* \neq *OO* *tr2*

\longrightarrow

($\exists sv1\ trv1\ sv2\ trv2.$

istateV *sv1* \wedge *istateV* *sv2* \wedge *corrState* *sv1* *s1* \wedge *corrState* *sv2* *s2* \wedge
Van.validFromS *sv1* *trv1* \wedge *Van.completedFrom* *sv1* *trv1* \wedge
Van.validFromS *sv2* *trv2* \wedge *Van.completedFrom* *sv2* *trv2* \wedge
SV *trv1* = *SO* *tr1* \wedge *SV* *trv2* = *SO* *tr2* \wedge
AV *trv1* = *AV* *trv2* \wedge *OV* *trv1* \neq *OV* *trv2*))

unfolding *rsecure-def*

unfolding *Van.leakVia-def* *Opt.leakVia-def*

by *auto metis*

end

locale *Statewise-Attacker-Mod* = *System-Mod* *istate* *validTrans* *final*

for *istate* :: 'state \Rightarrow bool **and** *validTrans* :: 'state \times 'state \Rightarrow bool **and** *final* ::
'state \Rightarrow bool

+

fixes

isSec :: 'state \Rightarrow bool **and** *getSec* :: 'state \Rightarrow 'secret

and

isInt :: 'state \Rightarrow bool **and** *getInt* :: 'state \Rightarrow 'act \times 'obs

assumes *final-not-isInt*: $\bigwedge s. \text{final } s \Longrightarrow \neg \text{isInt } s$

and *final-not-isSec*: $\bigwedge s. \text{final } s \Longrightarrow \neg \text{isSec } s$

begin

definition *getAct* :: 'state \Rightarrow 'act **where**

getAct = *fst* o *getInt*

definition *getObs* :: 'state \Rightarrow 'obs **where**

getObs = *snd* o *getInt*

definition $eqObs\ trn1\ trn2 \equiv$
 $(isInt\ trn1 \longleftrightarrow isInt\ trn2) \wedge (isInt\ trn1 \longrightarrow getObs\ trn1 = getObs\ trn2)$

definition $eqAct\ trn1\ trn2 \equiv$
 $(isInt\ trn1 \longleftrightarrow isInt\ trn2) \wedge (isInt\ trn1 \longrightarrow getAct\ trn1 = getAct\ trn2)$

definition $A :: 'state\ trace \Rightarrow 'act\ list$ **where**
 $A\ tr \equiv filtermap\ isInt\ getAct\ (butlast\ tr)$

sublocale $A: FiltermapBL\ isInt\ getAct\ A$
apply *standard unfolding* $A-def$ **..**

definition $O :: 'state\ trace \Rightarrow 'obs\ list$ **where**
 $O\ tr \equiv filtermap\ isInt\ getObs\ (butlast\ tr)$

sublocale $O: FiltermapBL\ isInt\ getObs\ O$
apply *standard unfolding* $O-def$ **..**

definition $S :: 'state\ list \Rightarrow 'secret\ list$ **where**
 $S\ tr \equiv filtermap\ isSec\ getSec\ (butlast\ tr)$

sublocale $S: FiltermapBL\ isSec\ getSec\ S$
apply *standard unfolding* $S-def$ **..**

end

sublocale $Statewise-Attacker-Mod < Attacker-Mod-fin$
where $S = S$ **and** $A = A$ **and** $O = O$
by *standard*

locale $Rel-Sec =$

$Van: Statewise-Attacker-Mod\ istateV\ validTransV\ finalV\ isSecV\ getSecV\ isIntV$
 $getIntV$

+

$Opt: Statewise-Attacker-Mod\ istateO\ validTransO\ finalO\ isSecO\ getSecO\ isIntO$
 $getIntO$

for $validTransV :: 'stateV \times 'stateV \Rightarrow bool$

and $istateV :: 'stateV \Rightarrow bool$ **and** $finalV :: 'stateV \Rightarrow bool$

and $isSecV :: 'stateV \Rightarrow bool$ **and** $getSecV :: 'stateV \Rightarrow 'secret$

and $isIntV :: 'stateV \Rightarrow bool$ **and** $getIntV :: 'stateV \Rightarrow 'actV \times 'obsV$

and $validTransO :: 'stateO \times 'stateO \Rightarrow bool$
and $istateO :: 'stateO \Rightarrow bool$ **and** $finalO :: 'stateO \Rightarrow bool$
and $isSecO :: 'stateO \Rightarrow bool$ **and** $getSecO :: 'stateO \Rightarrow 'secret$
and $isIntO :: 'stateO \Rightarrow bool$ **and** $getIntO :: 'stateO \Rightarrow 'actO \times 'obsO$

and $corrState :: 'stateV \Rightarrow 'stateO \Rightarrow bool$

sublocale $Rel\text{-}Sec < Relative\text{-}Security'\text{-}fin$
where $SV = Van.S$ **and** $AV = Van.A$ **and** $OV = Van.O$
and $SO = Opt.S$ **and** $AO = Opt.A$ **and** $OO = Opt.O$
by *standard*

context $Rel\text{-}Sec$
begin

abbreviation $getObsV :: 'stateV \Rightarrow 'obsV$ **where** $getObsV \equiv Van.getObs$
abbreviation $getActV :: 'stateV \Rightarrow 'actV$ **where** $getActV \equiv Van.getAct$
abbreviation $getObsO :: 'stateO \Rightarrow 'obsO$ **where** $getObsO \equiv Opt.getObs$
abbreviation $getActO :: 'stateO \Rightarrow 'actO$ **where** $getActO \equiv Opt.getAct$

abbreviation $reachV$ **where** $reachV \equiv Van.reach$
abbreviation $reachO$ **where** $reachO \equiv Opt.reach$

abbreviation $completedFromV :: 'stateV \Rightarrow 'stateV\ list \Rightarrow bool$ **where** $completedFromV \equiv Van.completedFrom$
abbreviation $completedFromO :: 'stateO \Rightarrow 'stateO\ list \Rightarrow bool$ **where** $completedFromO \equiv Opt.completedFrom$

lemmas $completedFromV\text{-}def = Van.completedFrom\text{-}def$
lemmas $completedFromO\text{-}def = Opt.completedFrom\text{-}def$

lemma $rsecure\text{-}def3$:

$rsecure \longleftrightarrow$

$(\forall s1\ tr1\ s2\ tr2.$

$istateO\ s1 \wedge Opt.validFromS\ s1\ tr1 \wedge completedFromO\ s1\ tr1 \wedge$
 $istateO\ s2 \wedge Opt.validFromS\ s2\ tr2 \wedge completedFromO\ s2\ tr2 \wedge$
 $Opt.A\ tr1 = Opt.A\ tr2 \wedge Opt.O\ tr1 \neq Opt.O\ tr2 \wedge$
 $(isIntO\ s1 \wedge isIntO\ s2 \longrightarrow getActO\ s1 = getActO\ s2)$

\longrightarrow

$(\exists sv1\ trv1\ sv2\ trv2.$

$istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge completedFromV\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge completedFromV\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge Van.O\ trv1 \neq Van.O\ trv2))$

unfolding $rsecure\text{-}def2$ **apply** (*intro iff-allI iffI impI*)

subgoal *by auto*

subgoal
by *clarsimp (metis (full-types) Opt.A.Cons-unfold
Opt.completed-Cons Opt.final-not-isInt
Simple-Transition-System.validFromS-Cons-iff
completedFromO-def list.sel(1) neq-Nil-conv) .*

definition *eqSec trnO trnA* \equiv
(isSecV trnO = isSecO trnA) \wedge (isSecV trnO \longrightarrow getSecV trnO = getSecO trnA)

lemma *eqSec-S-Cons'*:
eqSec trnO trnA \implies
(Van.S (trnO # trO') = Opt.S (trnA # trA')) \implies Van.S trO' = Opt.S trA'
apply(*cases trO' = []*)
subgoal apply(*cases trA' = []*)
subgoal by *auto*
subgoal unfolding *eqSec-def* **by** *auto .*
subgoal apply(*cases trA' = []*)
subgoal by *auto*
subgoal unfolding *eqSec-def* **by** *auto . .*

lemma *eqSec-S-Cons[simp]*:
eqSec trnO trnA \implies trO' = [] \longleftrightarrow trA' = [] \implies
(Van.S (trnO # trO') = Opt.S (trnA # trA')) \longleftrightarrow (Van.S trO' = Opt.S trA')
apply(*cases trO' = []*)
subgoal apply(*cases trA' = []*)
subgoal by *auto*
subgoal unfolding *eqSec-def* **by** *auto .*
subgoal apply(*cases trA' = []*)
subgoal by *auto*
subgoal unfolding *eqSec-def* **by** *auto . .*

end

locale *Relative-Security-Determ =*
Rel-Sec
validTransV istateV finalV isSecV getSecV isIntV getIntV
validTransO istateO finalO isSecO getSecO isIntO getIntO
corrState
+
System-Mod-Deterministic istateV validTransV finalV nextO
for *validTransV :: 'stateV \times 'stateV \Rightarrow bool*
and *istateV :: 'stateV \Rightarrow bool*
and *finalV :: 'stateV \Rightarrow bool*

```

and nextO :: 'stateV ⇒ 'stateV
and isSecV :: 'stateV ⇒ bool and getSecV :: 'stateV ⇒ 'secret
and isIntV :: 'stateV ⇒ bool and getIntV :: 'stateV ⇒ 'actV × 'obsV
and validTransO :: 'stateO × 'stateO ⇒ bool
and istateO :: 'stateO ⇒ bool
and finalO :: 'stateO ⇒ bool
and isSecO :: 'stateO ⇒ bool and getSecO :: 'stateO ⇒ 'secret
and isIntO :: 'stateO ⇒ bool and getIntO :: 'stateO ⇒ 'actO × 'obsO
and corrState :: 'stateV ⇒ 'stateO ⇒ bool

```

end

2 Relative Security

This theory formalizes the general notion of relative security, applicable to possibly infinite traces.

```

theory Relative-Security
imports Relative-Security-fin Preliminaries/Trivial
begin

```

```

no-notation relcomp (infixr O 75)
no-notation relcompp (infixr OO 75)

```

2.1 Leakage models and attacker models

```

locale Leakage-Mod = System-Mod istate validTrans final
for istate :: 'state ⇒ bool and validTrans :: 'state × 'state ⇒ bool and final ::
'state ⇒ bool
+
fixes leakVia :: 'state llist ⇒ 'state llist ⇒ 'leak ⇒ bool

```

```

locale Attacker-Mod = System-Mod istate validTrans final
for istate :: 'state ⇒ bool and validTrans :: 'state × 'state ⇒ bool and final ::
'state ⇒ bool
+
fixes S :: 'state llist ⇒ 'secret llist
and A :: 'state ltrace ⇒ 'act llist
and O :: 'state ltrace ⇒ 'obs llist
begin

```

```

fun leakVia :: 'state llist ⇒ 'state llist ⇒ 'secret llist × 'secret llist ⇒ bool
where
leakVia tr tr' (sl,sl') = (S tr = sl ∧ S tr' = sl' ∧ A tr = A tr' ∧ O tr ≠ O tr')

```

```

lemmas leakVia-def = leakVia.simps

```


end

sublocale *Attacker-Mod* < *Leakage-Mod*
where *leakVia* = *leakVia*
by *standard*

2.2 Locales for increasingly concrete notions of relative security

locale *Relative-Security''* =
 Van: *Leakage-Mod* *istateV* *validTransV* *finalV* *leakViaV*
 +
 Opt: *Leakage-Mod* *istateO* *validTransO* *finalO* *leakViaO*
 for *validTransV* :: 'stateV × 'stateV ⇒ bool
 and *istateV* :: 'stateV ⇒ bool **and** *finalV* :: 'stateV ⇒ bool
 and *leakViaV* :: 'stateV llist ⇒ 'stateV llist ⇒ 'leak ⇒ bool

 and *validTransO* :: 'stateO × 'stateO ⇒ bool
 and *istateO* :: 'stateO ⇒ bool **and** *finalO* :: 'stateO ⇒ bool
 and *leakViaO* :: 'stateO llist ⇒ 'stateO llist ⇒ 'leak ⇒ bool

 and *corrState* :: 'stateV ⇒ 'stateO ⇒ bool
begin

definition *lrsecure* :: bool **where**

lrsecure ≡ ∀ l s1 tr1 s2 tr2.
 istateO s1 ∧ *Opt.linvalidFromS* s1 tr1 ∧ *Opt.lcompletedFrom* s1 tr1 ∧
 istateO s2 ∧ *Opt.linvalidFromS* s2 tr2 ∧ *Opt.lcompletedFrom* s2 tr2 ∧
 leakViaO tr1 tr2 l
 →
 (∃ sv1 trv1 sv2 trv2.
 istateV sv1 ∧ *istateV* sv2 ∧ *corrState* sv1 s1 ∧ *corrState* sv2 s2 ∧
 Van.linvalidFromS sv1 trv1 ∧ *Van.lcompletedFrom* sv1 trv1 ∧
 Van.linvalidFromS sv2 trv2 ∧ *Van.lcompletedFrom* sv2 trv2 ∧
 leakViaV trv1 trv2 l)

end

locale *Relative-Security'* =
 Van: *Attacker-Mod* *istateV* *validTransV* *finalV* *SV* *AV* *OV*
 +
 Opt: *Attacker-Mod* *istateO* *validTransO* *finalO* *SO* *AO* *OO*
 for *validTransV* :: 'stateV × 'stateV ⇒ bool
 and *istateV* :: 'stateV ⇒ bool **and** *finalV* :: 'stateV ⇒ bool
 and *SV* :: 'stateV llist ⇒ 'secret llist

```

and AV :: 'stateV ltrace ⇒ 'actV llist
and OV :: 'stateV ltrace ⇒ 'obsV llist

and validTransO :: 'stateO × 'stateO ⇒ bool
and istateO :: 'stateO ⇒ bool and finalO :: 'stateO ⇒ bool
and SO :: 'stateO llist ⇒ 'secret llist
and AO :: 'stateO ltrace ⇒ 'actO llist
and OO :: 'stateO ltrace ⇒ 'obsO llist
and corrState :: 'stateV ⇒ 'stateO ⇒ bool

sublocale Relative-Security' < Relative-Security''
where leakViaV = Van.lleakVia and leakViaO = Opt.lleakVia
by standard

```

```

context Relative-Security'
begin

```

lemma lrsecure-def2:

```

lrsecure ↔
(∀ s1 tr1 s2 tr2.
  istateO s1 ∧ Opt.lvalidFromS s1 tr1 ∧ Opt.lcompletedFrom s1 tr1 ∧
  istateO s2 ∧ Opt.lvalidFromS s2 tr2 ∧ Opt.lcompletedFrom s2 tr2 ∧
  AO tr1 = AO tr2 ∧ OO tr1 ≠ OO tr2
  →
  (∃ sv1 trv1 sv2 trv2.
    istateV sv1 ∧ istateV sv2 ∧ corrState sv1 s1 ∧ corrState sv2 s2 ∧
    Van.lvalidFromS sv1 trv1 ∧ Van.lcompletedFrom sv1 trv1 ∧
    Van.lvalidFromS sv2 trv2 ∧ Van.lcompletedFrom sv2 trv2 ∧
    SV trv1 = SO tr1 ∧ SV trv2 = SO tr2 ∧
    AV trv1 = AV trv2 ∧ OV trv1 ≠ OV trv2))

```

```

unfolding lrsecure-def
unfolding Van.lleakVia-def Opt.lleakVia-def
by auto metis

```

end

```

context Statewise-Attacker-Mod begin

```

```

definition lA :: 'state ltrace ⇒ 'act llist where
lA tr ≡ lfiltermap isInt getAct (lbutlast tr)

```

sublocale *lA*: *LfiltermapBL isInt getAct lA*
apply standard unfolding *lA-def ..*

lemma *lA*: *lcompletedFrom s tr \implies lA tr = lmap getAct (lfilter isInt tr)*
apply(*cases lfinite tr*)
subgoal unfolding *lA.lmap-lfilter lbutlast-def*
by simp (*metis final-not-isInt lbutlast-lfinite lcompletedFrom-def lfilter-llist-of lfiltermap-lmap-lfilter lfinite-lfiltermap-butlast llast-llist-of llist-of-list-of lmap-llist-of*)
subgoal unfolding *lA.lmap-lfilter lbutlast-def* **by auto** .

definition *lO* :: *'state ltrace \Rightarrow 'obs llist* **where**
lO tr \equiv lfiltermap isInt getObs (lbutlast tr)

sublocale *lO*: *LfiltermapBL isInt getObs lO*
apply standard unfolding *lO-def ..*

lemma *lO*: *lcompletedFrom s tr \implies lO tr = lmap getObs (lfilter isInt tr)*
apply(*cases lfinite tr*)
subgoal unfolding *lO.lmap-lfilter lbutlast-def*
by simp (*metis List-Filtermap.filtermap-def butlast.simps(1) filtermap-butlast final-not-isInt lcompletedFrom-def lfilter-llist-of llist-of-list-of lmap-llist-of*)
subgoal unfolding *lO.lmap-lfilter lbutlast-def* **by auto** .

definition *lS* :: *'state llist \Rightarrow 'secret llist* **where**
lS tr \equiv lfiltermap isSec getSec (lbutlast tr)

sublocale *lS*: *LfiltermapBL isSec getSec lS*
apply standard unfolding *lS-def ..*

lemma *lS*: *lcompletedFrom s tr \implies lS tr = lmap getSec (lfilter isSec tr)*
apply(*cases lfinite tr*)
subgoal unfolding *lS.lmap-lfilter lbutlast-def*
by simp (*metis List-Filtermap.filtermap-def filtermap-butlast final-not-isSec lcompletedFrom-def lfilter-llist-of llist-of-eq-LNil-conv llist-of-list-of lmap-llist-of*)
subgoal unfolding *lS.lmap-lfilter lbutlast-def* **by auto** .

end

sublocale *Statewise-Attacker-Mod < Attacker-Mod*
where *S = lS* **and** *A = lA* **and** *O = lO*
by standard

sublocale *Rel-Sec* < *Relative-Security'*
where *SV* = *Van.lS* **and** *AV* = *Van.lA* **and** *OV* = *Van.lO*
and *SO* = *Opt.lS* **and** *AO* = *Opt.lA* **and** *OO* = *Opt.lO*
by *standard*

context *Rel-Sec*
begin

abbreviation *lcompletedFromV* :: 'stateV ⇒ 'stateV llist ⇒ bool **where** *lcompletedFromV* ≡ *Van.lcompletedFrom*
abbreviation *lcompletedFromO* :: 'stateO ⇒ 'stateO llist ⇒ bool **where** *lcompletedFromO* ≡ *Opt.lcompletedFrom*

lemma *eqSec-lS-Cons'*:
 $eqSec\ trnO\ trnA \implies$
 $(Van.lS\ (trnO\ \$\ trO') = Opt.lS\ (trnA\ \$\ trA')) \implies Van.lS\ trO' = Opt.lS\ trA'$
apply(cases *trO'* = [])
 subgoal apply(cases *trA'* = [])
 subgoal by *auto*
 subgoal unfolding *eqSec-def* **by** *auto* .
 subgoal apply(cases *trA'* = [])
 subgoal by *auto*
 subgoal unfolding *eqSec-def* **by** *auto* . .

lemma *eqSec-lS-Cons[simp]*:
 $eqSec\ trnO\ trnA \implies trO' = [] \longleftrightarrow trA' = [] \implies$
 $(Van.lS\ (trnO\ \$\ trO') = Opt.lS\ (trnA\ \$\ trA')) \longleftrightarrow (Van.lS\ trO' = Opt.lS\ trA')$
apply(cases *trO'* = [])
 subgoal apply(cases *trA'* = [])
 subgoal by *auto*
 subgoal unfolding *eqSec-def* **by** *auto* .
 subgoal apply(cases *trA'* = [])
 subgoal by *auto*
 subgoal unfolding *eqSec-def* **by** *auto* . .

end

end

3 Unwinding Proof Method for Finitary Relative Security

This theory formalizes the notion of unwinding for finitary relative security, and proves its soundness.

```

theory Unwinding-fin
imports Relative-Security
begin

```

3.1 The types and operators underlying unwinding: status, matching operators, etc.

```

context Rel-Sec
begin

```

```

datatype status = Eq | Diff

```

```

fun newStat :: status  $\Rightarrow$  bool  $\times$  'a  $\Rightarrow$  bool  $\times$  'a  $\Rightarrow$  status where
  newStat Eq (True,a) (True,a') = (if a = a' then Eq else Diff)
| newStat stat - - = stat

```

```

definition sstatO' statO sv1 sv2 = newStat statO (isIntV sv1, getObsV sv1)
(isIntV sv2, getObsV sv2)

```

```

definition sstatA' statA s1 s2 = newStat statA (isIntO s1, getObsO s1) (isIntO
s2, getObsO s2)

```

```

definition initCond ::

```

```

(enat  $\Rightarrow$  'stateO  $\Rightarrow$  'stateO  $\Rightarrow$  status  $\Rightarrow$  'stateV  $\Rightarrow$  'stateV  $\Rightarrow$  status  $\Rightarrow$  bool)  $\Rightarrow$ 
bool where

```

```

initCond  $\Delta \equiv \forall s1 s2.$ 

```

```

  istateO s1  $\wedge$  istateO s2

```

```

 $\longrightarrow$ 

```

```

  ( $\exists sv1 sv2. istateV sv1 \wedge istateV sv2 \wedge corrState sv1 s1 \wedge corrState sv2 s2$ 
 $\wedge \Delta \infty s1 s2 Eq sv1 sv2 Eq$ )

```

```

definition match1-1  $\Delta s1 s1' s2 statA sv1 sv2 statO \equiv$ 

```

```

 $\exists sv1'. validTransV (sv1,sv1') \wedge$ 
 $\Delta \infty s1' s2 statA sv1' sv2 statO$ 

```

```

definition match1-12  $\Delta s1 s1' s2 statA sv1 sv2 statO \equiv$ 

```

```

 $\exists sv1' sv2'.$ 

```

```

  let statO' = sstatO' statO sv1 sv2 in

```

```

  validTransV (sv1,sv1')  $\wedge$ 

```

```

  validTransV (sv2,sv2')  $\wedge$ 

```

```

 $\Delta \infty s1' s2 statA sv1' sv2' statO'$ 

```

```

definition match1  $\Delta s1 s2 statA sv1 sv2 statO \equiv$ 

```

```

 $\neg isIntO s1 \longrightarrow$ 

```

```

  ( $\forall s1'. validTransO (s1,s1')$ 

```

```

 $\longrightarrow$ 

```

$$\begin{aligned}
& (\neg \text{isSecO } s1 \wedge \Delta \infty s1' s2 \text{ statA } sv1 sv2 \text{ statO}) \vee \\
& (\text{eqSec } sv1 s1 \wedge \neg \text{isIntV } sv1 \wedge \text{match1-1 } \Delta s1 s1' s2 \text{ statA } sv1 sv2 \text{ statO}) \vee \\
& (\text{eqSec } sv1 s1 \wedge \neg \text{isSecV } sv2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match1-12 } \Delta s1 s1' s2 \\
& \text{statA } sv1 sv2 \text{ statO}))
\end{aligned}$$

lemmas *match1-defs* = *match1-def match1-1-def match1-12-def*

lemma *match1-1-mono*:

$$\Delta \leq \Delta' \implies \text{match1-1 } \Delta s1 s1' s2 \text{ statA } sv1 sv2 \text{ statO} \implies \text{match1-1 } \Delta' s1 s1' s2 \text{ statA } sv1 sv2 \text{ statO}$$

unfolding *le-fun-def match1-1-def* **by** *auto*

lemma *match1-12-mono*:

$$\Delta \leq \Delta' \implies \text{match1-12 } \Delta s1 s1' s2 \text{ statA } sv1 sv2 \text{ statO} \implies \text{match1-12 } \Delta' s1 s1' s2 \text{ statA } sv1 sv2 \text{ statO}$$

unfolding *le-fun-def match1-12-def* **by** *fastforce*

lemma *match1-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{match1 } \Delta s1 s2 \text{ statA } sv1 sv2 \text{ statO} \implies \text{match1 } \Delta' s1 s2 \text{ statA } sv1 sv2 \text{ statO}$

unfolding *match1-def* **apply** *clarify subgoal for s1'* **apply**(*erule alle[of - s1']*)

using *match1-1-mono[OF assms, of s1 s1' s2 statA sv1 sv2 statO]*

match1-12-mono[OF assms, of s1 s1' s2 statA sv1 sv2 statO]

assms[unfolded le-fun-def, rule-format, of - s1' s2 statA sv1 sv2 statO]

by *auto* .

definition *match2-1* $\Delta s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO} \equiv$

$$\begin{aligned}
& \exists sv2'. \text{validTransV } (sv2, sv2') \wedge \\
& \Delta \infty s1 s2' \text{ statA } sv1 sv2' \text{ statO}
\end{aligned}$$

definition *match2-12* $\Delta s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO} \equiv$

$$\begin{aligned}
& \exists sv1' sv2'. \\
& \text{let } \text{statO}' = \text{sstatO}' \text{ statO } sv1 sv2 \text{ in} \\
& \text{validTransV } (sv1, sv1') \wedge \\
& \text{validTransV } (sv2, sv2') \wedge \\
& \Delta \infty s1 s2' \text{ statA } sv1' sv2' \text{ statO}'
\end{aligned}$$

definition *match2* $\Delta s1 s2 \text{ statA } sv1 sv2 \text{ statO} \equiv$

$$\begin{aligned}
& \neg \text{isIntO } s2 \longrightarrow \\
& (\forall s2'. \text{validTransO } (s2, s2') \\
& \longrightarrow \\
& (\neg \text{isSecO } s2 \wedge \Delta \infty s1 s2' \text{ statA } sv1 sv2 \text{ statO}) \vee \\
& (\text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge \text{match2-1 } \Delta s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO}) \vee \\
& (\neg \text{isSecV } sv1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match2-12 } \Delta s1 s2 s2' \\
& \text{statA } sv1 sv2 \text{ statO}))
\end{aligned}$$

lemmas *match2-defs* = *match2-def match2-1-def match2-12-def*

lemma *match2-1-mono*:

$\Delta \leq \Delta' \implies \text{match2-1 } \Delta \ s1 \ s1' \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \implies \text{match2-1 } \Delta' \ s1 \ s1' \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$

unfolding *le-fun-def match2-1-def* **by** *auto*

lemma *match2-12-mono*:

$\Delta \leq \Delta' \implies \text{match2-12 } \Delta \ s1 \ s1' \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \implies \text{match2-12 } \Delta' \ s1 \ s1' \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$

unfolding *le-fun-def match2-12-def* **by** *fastforce*

lemma *match2-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{match2 } \Delta \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \implies \text{match2 } \Delta' \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$

unfolding *match2-def* **apply** *clarify subgoal for s2'* **apply** (*erule* *allE*[*of* - *s2'*])

using *match2-1-mono*[*OF* *assms*, *of* *s1 s2 s2' statA sv1 sv2 statO*]

match2-12-mono[*OF* *assms*, *of* *s1 s2 s2' statA sv1 sv2 statO*]

assms[*unfolded le-fun-def*, *rule-format*, *of* - *s1 s2' statA sv1 sv2 statO*]

by *auto* .

definition *match12-1* $\Delta \ s1' \ s2' \ \text{statA}' \ sv1 \ sv2 \ \text{statO} \equiv$

$\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 $\Delta \infty \ s1' \ s2' \ \text{statA}' \ sv1' \ sv2 \ \text{statO}$

definition *match12-2* $\Delta \ s1' \ s2' \ \text{statA}' \ sv1 \ sv2 \ \text{statO} \equiv$

$\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty \ s1' \ s2' \ \text{statA}' \ sv1 \ sv2' \ \text{statO}$

definition *match12-12* $\Delta \ s1' \ s2' \ \text{statA}' \ sv1 \ sv2 \ \text{statO} \equiv$

$\exists sv1' \ sv2'.$
let *statO'* = *sstatO'* *statO* *sv1 sv2* *in*
 $\text{validTransV } (sv1, sv1') \wedge$
 $\text{validTransV } (sv2, sv2') \wedge$
 $(\text{statA}' = \text{Diff} \longrightarrow \text{statO}' = \text{Diff}) \wedge$
 $\Delta \infty \ s1' \ s2' \ \text{statA}' \ sv1' \ sv2' \ \text{statO}'$

definition *match12* $\Delta \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \equiv$

$\forall s1' \ s2'.$

let *statA'* = *sstatA'* *statA* *s1 s2* *in*
 $\text{validTransO } (s1, s1') \wedge$
 $\text{validTransO } (s2, s2') \wedge$
 $\text{Opt.eqAct } s1 \ s2 \wedge$
 $\text{isIntO } s1 \wedge \text{isIntO } s2$

\longrightarrow

$(\neg \text{isSecO } s1 \wedge \neg \text{isSecO } s2 \wedge (\text{statA} = \text{statA}' \vee \text{statO} = \text{Diff}) \wedge \Delta \infty \ s1' \ s2')$

$statA' sv1 sv2 statO$
 \vee
 $(\neg isSecO s2 \wedge eqSec sv1 s1 \wedge \neg isIntV sv1 \wedge$
 $(statA = statA' \vee statO = Diff) \wedge$
 $match12-1 \Delta s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\neg isSecO s1 \wedge eqSec sv2 s2 \wedge \neg isIntV sv2 \wedge$
 $(statA = statA' \vee statO = Diff) \wedge$
 $match12-2 \Delta s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(eqSec sv1 s1 \wedge eqSec sv2 s2 \wedge Van.eqAct sv1 sv2 \wedge$
 $match12-12 \Delta s1' s2' statA' sv1 sv2 statO)$

lemmas $match12-defs = match12-def match12-1-def match12-2-def match12-12-def$

lemma $match12-simpleI$:

assumes $\bigwedge s1' s2' statA'$.

$statA' = sstatA' statA s1 s2 \implies$

$validTransO (s1, s1') \implies$

$validTransO (s2, s2') \implies$

$Opt.eqAct s1 s2 \implies$

$(\neg isSecO s1 \wedge \neg isSecO s2 \wedge (statA = statA' \vee statO = Diff) \wedge \Delta \infty s1' s2'$
 $statA' sv1 sv2 statO)$

\vee

$(eqSec sv1 s1 \wedge eqSec sv2 s2 \wedge Van.eqAct sv1 sv2 \wedge$

$match12-12 \Delta s1' s2' statA' sv1 sv2 statO)$

shows $match12 \Delta s1 s2 statA sv1 sv2 statO$

using $assms unfolding match12-def Let-def$ **by** $blast$

lemma $match12-1-mono$:

$\Delta \leq \Delta' \implies match12-1 \Delta s1' s2' statA' sv1 sv2 statO \implies match12-1 \Delta' s1' s2'$
 $statA' sv1 sv2 statO$

unfolding $le-fun-def match12-1-def$ **by** $auto$

lemma $match12-2-mono$:

$\Delta \leq \Delta' \implies match12-2 \Delta s1 s2' statA' sv1 sv2 statO \implies match12-2 \Delta' s1 s2'$
 $statA' sv1 sv2 statO$

unfolding $le-fun-def match12-2-def$ **by** $auto$

lemma $match12-12-mono$:

$\Delta \leq \Delta' \implies match12-12 \Delta s1' s2' statA' sv1 sv2 statO \implies match12-12 \Delta' s1'$
 $s2' statA' sv1 sv2 statO$

unfolding $le-fun-def match12-12-def$ **by** $fastforce$

lemma $match12-mono$:

assumes $\Delta \leq \Delta'$

shows $match12 \Delta s1 s2 statA sv1 sv2 statO \implies match12 \Delta' s1 s2 statA sv1 sv2$
 $statO$

unfolding *match12-def* **apply** *clarify subgoal for* $s1' s2'$ **apply**(*erule allE*[of - $s1'$]) **apply**(*erule allE*[of - $s2'$])
using *match12-1-mono*[OF *assms*, of $s1' s2' - sv1 sv2 statO$]
match12-2-mono[OF *assms*, of $s1' s2' - sv1 sv2 statO$]
match12-12-mono[OF *assms*, of $s1' s2' - sv1 sv2 statO$]
assms[*unfolded le-fun-def*, *rule-format*, of - $s1' s2'$
sstatA' statA s1 s2 sv1 sv2 statO]
by *simp metis* .

definition *react* $\Delta s1 s2 statA sv1 sv2 statO \equiv$
match1 $\Delta s1 s2 statA sv1 sv2 statO$
 \wedge
match2 $\Delta s1 s2 statA sv1 sv2 statO$
 \wedge
match12 $\Delta s1 s2 statA sv1 sv2 statO$

lemmas *react-defs* = *match1-def match2-def match12-def*
lemmas *match-deep-defs* = *match1-defs match2-defs match12-defs*

lemma *match-mono*:

assumes $\Delta \leq \Delta'$

shows *react* $\Delta s1 s2 statA sv1 sv2 statO \implies$ *react* $\Delta' s1 s2 statA sv1 sv2 statO$

unfolding *react-def* **using** *match1-mono*[OF *assms*] *match2-mono*[OF *assms*] *match12-mono*[OF *assms*] **by** *auto*

definition *move-1* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 $\Delta w s1 s2 statA sv1' sv2 statO$

definition *move-2* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta w s1 s2 statA sv1 sv2' statO$

definition *move-12* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv1' sv2'.$
let *statO' = sstatO' statO sv1 sv2 in*
 $\text{validTransV } (sv1, sv1') \wedge \text{validTransV } (sv2, sv2') \wedge$
 $\Delta w s1 s2 statA sv1' sv2' statO'$

definition *proact* $\Delta w s1 s2 statA sv1 sv2 statO \equiv$
 $(\neg \text{isSecV } sv1 \wedge \neg \text{isIntV } sv1 \wedge \text{move-1 } \Delta w s1 s2 statA sv1 sv2 statO)$
 \vee
 $(\neg \text{isSecV } sv2 \wedge \neg \text{isIntV } sv2 \wedge \text{move-2 } \Delta w s1 s2 statA sv1 sv2 statO)$
 \vee
 $(\neg \text{isSecV } sv1 \wedge \neg \text{isSecV } sv2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{move-12 } \Delta w s1 s2 statA$

sv1 sv2 statO)

lemmas *proact-defs = proact-def move-1-def move-2-def move-12-def*

lemma *move-1-mono:*

$\Delta \leq \Delta' \implies \text{move-1 } \Delta \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO} \implies \text{move-1 } \Delta' \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

unfolding *le-fun-def move-1-def* **by** *auto*

lemma *move-2-mono:*

$\Delta \leq \Delta' \implies \text{move-2 } \Delta \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO} \implies \text{move-2 } \Delta' \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

unfolding *le-fun-def move-2-def* **by** *auto*

lemma *move-12-mono:*

$\Delta \leq \Delta' \implies \text{move-12 } \Delta \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO} \implies \text{move-12 } \Delta' \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

unfolding *le-fun-def move-12-def* **by** *fastforce*

lemma *proact-mono:*

assumes $\Delta \leq \Delta'$

shows $\text{proact } \Delta \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO} \implies \text{proact } \Delta' \text{ meas } s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

unfolding *proact-def* **using** *move-1-mono[OF assms] move-2-mono[OF assms] move-12-mono[OF assms]* **by** *auto*

3.2 The definition of unwinding

definition *unwindCond* ::

$(\text{enat} \implies 'stateO \implies 'stateO \implies \text{status} \implies 'stateV \implies 'stateV \implies \text{status} \implies \text{bool}) \implies \text{bool}$

where

$\text{unwindCond } \Delta \equiv \forall w \ s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}.$

$\text{reachO } s1 \wedge \text{reachO } s2 \wedge \text{reachV } sv1 \wedge \text{reachV } sv2 \wedge$

$\Delta \ w \ s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

\longrightarrow

$(\text{finalO } s1 \longleftrightarrow \text{finalO } s2) \wedge (\text{finalV } sv1 \longleftrightarrow \text{finalO } s1) \wedge (\text{finalV } sv2 \longleftrightarrow \text{finalO } s2)$

\wedge

$(\text{statA} = \text{Eq} \longrightarrow (\text{isIntO } s1 \longleftrightarrow \text{isIntO } s2))$

\wedge

$(\exists v < w. \text{proact } \Delta \ v \ s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO})$

\vee

$\text{react } \Delta \ s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

)

lemma *unwindCond-simpleI*:

assumes

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

and

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \implies statA = Eq$
 \implies
 $isIntO\ s1 \longleftrightarrow isIntO\ s2$

and

$\bigwedge w\ s1\ s2\ statA\ sv1\ sv2\ statO.$
 $reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$
 $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$
 \implies
 $react\ \Delta\ s1\ s2\ statA\ sv1\ sv2\ statO$

shows *unwindCond* Δ

using *assms unfolding unwindCond-def* **by** *auto*

3.3 The soundness of unwinding

definition $\psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2 \equiv$

$trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $(finalV\ (lastt\ sv1\ trv1) \longleftrightarrow finalO\ (lastt\ s1\ tr1)) \wedge (finalV\ (lastt\ sv2\ trv2) \longleftrightarrow$
 $finalO\ (lastt\ s2\ tr2)) \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \wedge Opt.O\ tr1 \neq Opt.O\ tr2 \longrightarrow Van.O\ trv1 \neq Van.O\ trv2)$

lemma *ψ -completedFrom*: $completedFromO\ s1\ tr1 \implies completedFromO\ s2\ tr2 \implies$

$\psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2$
 $\implies completedFromV\ sv1\ trv1 \wedge completedFromV\ sv2\ trv2$

unfolding *ψ -def Opt.completedFrom-def Van.completedFrom-def lastt-def*
by *presburger*

lemma *completedFromO-lastt*: $completedFromO\ s1\ tr1 \implies finalO\ (lastt\ s1\ tr1)$

unfolding *Opt.completedFrom-def lastt-def* **by** *auto*

lemma *rsecure-strong*:

assumes

$\wedge s1\ tr1\ s2\ tr2.$

$istateO\ s1 \wedge Opt.validFromS\ s1\ tr1 \wedge completedFromO\ s1\ tr1 \wedge$

$istateO\ s2 \wedge Opt.validFromS\ s2\ tr2 \wedge completedFromO\ s2\ tr2 \wedge$

$Opt.A\ tr1 = Opt.A\ tr2 \wedge (isIntO\ s1 \wedge isIntO\ s2 \longrightarrow getActO\ s1 = getActO$

$s2)$

\implies

$\exists sv1\ trv1\ sv2\ trv2.$

$istateV\ sv1 \wedge istateV\ sv2 \wedge corrState\ sv1\ s1 \wedge corrState\ sv2\ s2 \wedge$

$\psi\ s1\ tr1\ s2\ tr2\ Eq\ sv1\ trv1\ sv2\ trv2$

shows *rsecure*

unfolding *rsecure-def3* **apply** *clarify*

subgoal for $s1\ tr1\ s2\ tr2$

using *assms*[*of* $s1\ tr1\ s2\ tr2$]

using ψ -*completedFrom* ψ -*def* *completedFromO-lastt* **by** (*metis* (*full-types*))

.

proposition *unwindCond-ex- ψ* :

assumes *unwind*: *unwindCond* Δ

and Δ : $\Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO$ **and** *stat*: ($statA = Diff \longrightarrow statO = Diff$)

and *v*: $Opt.validFromS\ s1\ tr1\ Opt.completedFrom\ s1\ tr1\ Opt.validFromS\ s2\ tr2$
 $Opt.completedFrom\ s2\ tr2$

and *tr14*: $Opt.A\ tr1 = Opt.A\ tr2$

and *r*: $reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$

shows $\exists trv1\ trv2. \psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2$

using *assms*(2-)

proof(*induction* $length\ tr1 + length\ tr2\ w$

arbitrary: $s1\ s2\ statA\ sv1\ sv2\ statO\ tr1\ tr2$ *rule*: *less2-induct'*)

case (*less* $w\ tr1\ tr2\ s1\ s2\ statA\ sv1\ sv2\ statO$)

note $ok = \langle statA = Diff \longrightarrow statO = Diff \rangle$

note $\Delta = \langle \Delta\ w\ s1\ s2\ statA\ sv1\ sv2\ statO \rangle$

note $A34 = \langle Opt.A\ tr1 = Opt.A\ tr2 \rangle$

note $r34 = less.prem(8,9)$ **note** $r12 = less.prem(10,11)$

note $r = r34\ r12$

note $r3 = r34(1)$ **note** $r4 = r34(2)$ **note** $r1 = r12(1)$ **note** $r2 = r12(2)$

have $i34$: $statA = Eq \longrightarrow isIntO\ s1 = isIntO\ s2$

and $f34$: $finalO\ s1 = finalO\ s2 \wedge finalV\ sv1 = finalO\ s1 \wedge finalV\ sv2 = finalO$
 $s2$

using Δ *unwind*[*unfolded* *unwindCond-def*] *r* **by** *auto*

have *proact-match*: ($\exists v < w. proact\ \Delta\ v\ s1\ s2\ statA\ sv1\ sv2\ statO$) $\vee react\ \Delta\ s1$
 $s2\ statA\ sv1\ sv2\ statO$

using Δ *unwind*[*unfolded* *unwindCond-def*] *r* **by** *auto*

show *?case* **using** *proact-match* **proof** *safe*

fix *v* **assume** *v*: $v < w$

assume *proact* $\Delta\ v\ s1\ s2\ statA\ sv1\ sv2\ statO$

```

thus ?thesis unfolding proact-def proof safe
  assume sv1:  $\neg$  isSecV sv1  $\neg$  isIntV sv1 and move-1  $\Delta$  v s1 s2 statA sv1 sv2
statO
  then obtain sv1'
  where 0: validTransV (sv1,sv1')
  and  $\Delta$ :  $\Delta$  v s1 s2 statA sv1' sv2 statO
  unfolding move-1-def by auto
  have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2 tr2 statO sv1' trv1 sv2 trv2
  using less(2)[OF v, of tr1 tr2 s1 s2 statA sv1' sv2 statO, simplified, OF  $\Delta$ 
ok - - - - r34 r1' r2]
  using A34 less.prem(3-6) by blast
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
  using  $\psi$  ok 0 sv1 unfolding  $\psi$ -def Van.completedFrom-def by auto
next
  assume sv2:  $\neg$  isSecV sv2  $\neg$  isIntV sv2 and move-2  $\Delta$  v s1 s2 statA sv1 sv2
statO
  then obtain sv2'
  where 0: validTransV (sv2,sv2')
  and  $\Delta$ :  $\Delta$  v s1 s2 statA sv1 sv2' statO
  unfolding move-2-def by auto
  have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2 tr2 statO sv1 trv1 sv2' trv2
  using less(2)[OF v, of tr1 tr2 s1 s2 statA sv1 sv2' statO, simplified, OF  $\Delta$ 
ok - - - - r34 r1 r2']
  using A34 less.prem(3-6) by blast
  show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\psi$  ok 0 sv2 unfolding  $\psi$ -def Van.completedFrom-def by auto
next
  assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
  and move-12  $\Delta$  v s1 s2 statA sv1 sv2 statO
  then obtain sv1' sv2' statO'
  where 0: statO' = sstatO' statO sv1 sv2
  validTransV (sv1,sv1')  $\neg$  isSecV sv1
  validTransV (sv2,sv2')  $\neg$  isSecV sv2
  Van.eqAct sv1 sv2
  and  $\Delta$ :  $\Delta$  v s1 s2 statA sv1' sv2' statO'
  unfolding move-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
  have ok': statA = Diff  $\longrightarrow$  statO' = Diff using ok 0 unfolding sstatO'-def
by (cases statO, auto)
  obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2 tr2 statO' sv1' trv1 sv2' trv2
  using less(2)[OF v, of tr1 tr2 s1 s2 statA sv1' sv2' statO', simplified, OF  $\Delta$ 
ok' - - - - r34 r12']
  using A34 less.prem(3-6) by blast
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
  using  $\psi$  ok' 0 sv12 unfolding  $\psi$ -def sstatO'-def Van.completedFrom-def

```

```

using Van.A.Cons-unfold Van.eqAct-def completedFromO-lastt less.prems(4)
less.prems(6) by auto
qed
next
assume m: react Δ s1 s2 statA sv1 sv2 statO
show ?thesis
proof(cases length tr1 ≤ Suc 0)
  case True note tr1 = True
  hence tr1 = [] ∨ tr1 = [s1]
  by (metis Opt.validFromS-Cons-iff le-0-eq le-SucE length-0-conv length-Suc-conv
less.prems(3))
    hence finalO s1 using less(3-6)
    using Opt.completed-Cons Opt.completed-Nil by blast
    hence f4: finalO s2 using f34 by blast
    hence tr2: tr2 = [] ∨ tr2 = [s2]
    by (metis Opt.final-def Simple-Transition-System.validFromS-Cons-iff less.prems(5)
neq-Nil-conv)
      show ?thesis apply(rule exI[of - [sv1]], rule exI[of - [sv2]]) using tr1 tr2
      using f4 f34
      using completedFromO-lastt less.prems(4)
      by (auto simp add: lastt-def ψ-def)
next
  case False
  then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1' ≠ []
    by (cases tr1, auto)
  have s13[simp]: s13 = s1 using ⟨Opt.validFromS s1 tr1⟩
    by (simp add: Opt.validFromS-Cons-iff tr1)
  obtain s1' where
    trn3: validTransO (s1, s1') and
    tr1': Opt.validFromS s1' tr1' using ⟨Opt.validFromS s1 tr1⟩
  unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
    have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv
snd-conv)
    have f3: ¬ finalO s1 using Opt.final-def trn3 by blast
    hence f4: ¬ finalO s2 using f34 by blast
    hence tr2: ¬ length tr2 ≤ Suc 0
    by (metis (no-types, opaque-lifting) Opt.completed-Cons Opt.completed-Nil
Simple-Transition-System.validFromS-Cons-iff Suc-n-not-le-n bot-nat-0.extremum
le-Suc-eq length-Cons less.prems(5) less.prems(6) list.exhaust order-antisym-conv)
    then obtain s24 tr2' where tr2: tr2 = s24 # tr2' and tr2'NE: tr2' ≠ []
    by (cases tr2, auto)
    have s24[simp]: s24 = s2 using ⟨Opt.validFromS s2 tr2⟩
    by (simp add: Opt.validFromS-Cons-iff tr2)
    obtain s2' where
      trn4: validTransO (s2, s2') ∨ (s2 = s2' ∧ tr2' = []) and
      tr2': Opt.validFromS s2' tr2' using ⟨Opt.validFromS s2 tr2⟩
    unfolding tr2 s24 using Opt.validFromS-Cons-iff by auto
    have r34': reachO s1' reachO s2'
    using r3 trn3 r4 trn4 by (metis Opt.reach.Step fst-conv snd-conv)+

```

```

note  $r3' = r34'(1)$  note  $r4' = r34'(2)$ 
define  $statA'$  where  $statA': statA' = sstatA' statA s1 s2$ 
have  $\neg isIntO s1 \vee \neg isIntO s2 \vee (isIntO s1 \wedge isIntO s2)$ 
by auto
thus ?thesis
proof safe
  assume  $isAO3: \neg isIntO s1$ 
  have  $O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'$ 
  using  $isAO3$  unfolding  $tr1$  by auto
  have  $A34': Opt.A tr1' = Opt.A tr2$ 
  using  $A34 O33'(2)$  by auto
  have  $m: match1 \Delta s1 s2 statA sv1 sv2 statO$  using  $m$  unfolding react-def
by auto
  have  $(\neg isSecO s1 \wedge \Delta \infty s1' s2 statA sv1 sv2 statO) \vee$ 
     $(eqSec sv1 s1 \wedge \neg isIntV sv1 \wedge match1-1 \Delta s1 s1' s2 statA sv1 sv2$ 
 $statO) \vee$ 
     $(eqSec sv1 s1 \wedge \neg isSecV sv2 \wedge Van.eqAct sv1 sv2 \wedge match1-12 \Delta s1$ 
 $s1' s2 statA sv1 sv2 statO)$ 
  using  $m isAO3 trn3 ok$  unfolding match1-def by auto
  thus ?thesis
  proof safe
    assume  $\neg isSecO s1$  and  $\Delta: \Delta \infty s1' s2 statA sv1 sv2 statO$ 
    hence  $S3: Opt.S tr1' = Opt.S tr1$  unfolding  $tr1$  by auto
    obtain  $trv1 trv2$  where  $\psi: \psi s1 tr1' s2 tr2 statO sv1 trv1 sv2 trv2$ 
    using  $less(1)[of tr1' tr2, OF - \Delta - - - - - r3' r4 r12, unfolded O33',$ 
simplified]
    using  $less.prem s tr1' ok A34' f3 f4$  unfolding  $tr1$  Opt.completedFrom-def
    by  $(auto split: if-splits simp: \psi-def lastt-def)$ 
    show ?thesis apply $(rule exI[of - trv1])$  apply $(rule exI[of - trv2])$ 
    using  $\psi O33' S3$  unfolding  $\psi-def$ 
    using completedFromO-lastt less.prem s(4)
    by  $(auto simp add: tr1 tr1' NE)$ 
  next
    assume  $trn13: eqSec sv1 s1$  and
     $Atrn1: \neg isIntV sv1$  and  $match1-1 \Delta s1 s1' s2 statA sv1 sv2 statO$ 
    then obtain  $sv1'$  where
     $trn1: validTransV (sv1, sv1')$  and
     $\Delta: \Delta \infty s1' s2 statA sv1' sv2 statO$ 
    unfolding match1-1-def by auto
    have  $r1': reachV sv1' using r1 trn1$  by  $(metis Van.reach.Step fst-conv$ 
snd-conv)
    obtain  $trv1 trv2$  where  $\psi: \psi s1 tr1' s2 tr2 statO sv1' trv1 sv2 trv2$ 
    using  $less(1)[of tr1' tr2, OF - \Delta - - - - - r3' r4 r1' r2, unfolded O33',$ 
simplified]
    using  $less.prem s tr1' ok A34' f3 f4$  unfolding  $tr1 tr2$  Opt.completedFrom-def

    by  $(auto simp: \psi-def lastt-def split: if-splits)$ 
    show ?thesis apply $(rule exI[of - sv1 \# trv1])$  apply $(rule exI[of - trv2])$ 
    using  $\psi O33'$  unfolding  $tr1 tr2$  Van.completedFrom-def

```

```

using Van.validFromS-Cons trn1 tr1'NE tr2'NE
using isAO3 ok Atrn1 eqSec-S-Cons trn13
unfolding  $\psi$ -def using completedFromO-lastt less.premis(4) tr1 by auto

next
  assume sv2:  $\neg$  isSecV sv2 and trn13: eqSec sv1 s1 and
  Atrn12: Van.eqAct sv1 sv2 and match1-12  $\Delta$  s1 s1' s2 statA sv1 sv2 statO
  then obtain sv1' sv2' statO' where
    statO': statO' = sstatO' statO sv1 sv2 and
    trn1: validTransV (sv1,sv1') and
    trn2: validTransV (sv2,sv2') and
     $\Delta$ :  $\Delta \infty$  s1' s2 statA sv1' sv2' statO'
  unfolding match1-12-def by auto
  have r12': reachV sv1' reachV sv2'
  using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
  obtain trv1 trv2 where  $\psi$ :  $\psi$  s1' tr1' s2 tr2 statO' sv1' trv1 sv2' trv2
    using less(1)[of tr1' tr2, OF -  $\Delta$  - - - - - r3' r4 r12', unfolded O33',
simplified]
  using less.premis tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
  by auto presburger+
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
  using  $\psi$  O33' tr1'NE tr2'NE sv2
  using Van.validFromS-Cons trn1 trn2
  using isAO3 ok Atrn12 eqSec-S-Cons trn13 f3 f34 s13
unfolding  $\psi$ -def tr1 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def
  using Van.A.Cons-unfold tr1' trn3 by auto
qed
next
  assume isAO4:  $\neg$  isIntO s2
  have O44': Opt.O tr2 = Opt.O tr2' Opt.A tr2 = Opt.A tr2'
  using isAO4 unfolding tr2 by auto
  have A34': Opt.A tr1 = Opt.A tr2'
  using A34 O44'(2) by auto
  have m: match2  $\Delta$  s1 s2 statA sv1 sv2 statO using m unfolding react-def
by auto
  have ( $\neg$  isSecO s2  $\wedge$   $\Delta \infty$  s1 s2' statA sv1 sv2 statO)  $\vee$ 
    (eqSec sv2 s2  $\wedge$   $\neg$  isIntV sv2  $\wedge$  match2-1  $\Delta$  s1 s2 s2' statA sv1 sv2
statO)  $\vee$ 
    ( $\neg$  isSecV sv1  $\wedge$  eqSec sv2 s2  $\wedge$  Van.eqAct sv1 sv2  $\wedge$  match2-12  $\Delta$  s1
s2 s2' statA sv1 sv2 statO)
  using m isAO4 trn4 ok tr2'NE unfolding match2-def by auto
  thus ?thesis
proof safe
  assume  $\neg$  isSecO s2 and  $\Delta$ :  $\Delta \infty$  s1 s2' statA sv1 sv2 statO
  hence S4: Opt.S tr2' = Opt.S tr2 unfolding tr2 by auto
  obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2' tr2' statO sv1 trv1 sv2 trv2
  using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4', simplified]

```



```

    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
  by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
    using  $\psi$  O44' S4 unfolding  $\psi$ -def
    using completedFromO-lastt less.premis(6)
    unfolding Opt.completedFrom-def using tr2 tr2'NE by auto
  next
    assume trn24: eqSec sv2 s2 and
      Atrn2:  $\neg$  isIntV sv2 and match2-1  $\Delta$  s1 s2 s2' statA sv1 sv2 statO
    then obtain sv2' where trn2: validTransV (sv2,sv2') and
       $\Delta$ :  $\Delta \infty$  s1 s2' statA sv1 sv2' statO
    unfolding match2-1-def by auto
    have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
    obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2' tr2' statO sv1 trv1 sv2' trv2
    using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4' r1 r2', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
  by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
    using  $\psi$  tr1'NE tr2'NE
    using Van.validFromS-Cons trn2
    using isAO4 ok Atrn2 eqSec-S-Cons trn24
    unfolding  $\psi$ -def tr1 tr2 s13 s24 Van.completedFrom-def lastt-def by auto
  next
    assume sv1:  $\neg$  isSecV sv1 and trn24: eqSec sv2 s2 and
      Atrn12: Van.eqAct sv1 sv2 and match2-12  $\Delta$  s1 s2 s2' statA sv1 sv2
  statO
    then obtain sv1' sv2' statO' where
      statO': statO' = sstatO' statO sv1 sv2 and
      trn1: validTransV (sv1,sv1') and
      trn2: validTransV (sv2,sv2') and
       $\Delta$ :  $\Delta \infty$  s1 s2' statA sv1' sv2' statO'
    unfolding match2-12-def by auto
    have r12': reachV sv1' reachV sv2'
    using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
    obtain trv1 trv2 where  $\psi$ :  $\psi$  s1 tr1 s2' tr2' statO' sv1' trv1 sv2' trv2
    using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4' r12', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
  statO' sstatO'-def
    by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
    using  $\psi$  O44' tr1'NE tr2'NE sv1
    using Van.validFromS-Cons trn1 trn2
    using isAO4 ok Atrn12 eqSec-S-Cons trn24
    unfolding  $\psi$ -def tr2 tr1'NE Van.completedFrom-def Van.eqAct-def
      statO' sstatO'-def
    using Van.A.Cons-unfold tr2' trn4 by auto
  qed

```

```

next
  assume isAO34: isIntO s1 isIntO s2
  have A34': getActO s1 = getActO s2 Opt.A tr1' = Opt.A tr2'
  using A34 isAO34 tr1'NE tr2'NE unfolding tr1 tr2 by auto
  have O33': Opt.O tr1 = getObsO s1 # Opt.O tr1' and
    O44': Opt.O tr2 = getObsO s2 # Opt.O tr2'
  using isAO34 tr1'NE tr2'NE unfolding tr1 s13 tr2 s24 by auto
  have m: match12 Δ s1 s2 statA sv1 sv2 statO using m unfolding statA'
  react-def by auto
  have trn34: getObsO s1 = getObsO s2 ∨ statA' = Diff
  using isAO34 unfolding statA' sstatA'-def by (cases statA, auto)
  have (¬ isSecO s1 ∧ ¬ isSecO s2 ∧ (statA = statA' ∨ statO = Diff) ∧ Δ
  ∞ s1' s2' statA' sv1 sv2 statO)
    ∨
    (¬ isSecO s2 ∧ eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧
    (statA = statA' ∨ statO = Diff) ∧
    match12-1 Δ s1' s2' statA' sv1 sv2 statO)
    ∨
    (¬ isSecO s1 ∧ eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧
    (statA = statA' ∨ statO = Diff) ∧
    match12-2 Δ s1' s2' statA' sv1 sv2 statO)
    ∨
    (eqSec sv1 s1 ∧ eqSec sv2 s2 ∧ Van.eqAct sv1 sv2 ∧
    match12-12 Δ s1' s2' statA' sv1 sv2 statO)
  (is ?K1 ∨ ?K2 ∨ ?K3 ∨ ?K4)
  using m[unfolded match12-def, rule-format, of s1' s2']
  isAO34 A34' trn3 trn4 tr1'NE tr2'NE
  unfolding s13 s24 trn34 statA' Opt.eqAct-def sstatA'-def by auto
  thus ?thesis proof (elim disjE)
    assume K1: ?K1 hence Δ: Δ ∞ s1' s2' statA' sv1 sv2 statO by simp
    have ok': (statA' = Diff → statO = Diff)
    using ok K1 unfolding statA' using isAO34 by auto
    obtain trv1 trv2 where ψ: ψ s1' tr1' s2' tr2' statO sv1 trv1 sv2 trv2
    using less(1)[of tr1' tr2', OF - Δ - - - - - r34' r12, simplified]
    using less.premis tr1' tr2' ok' A34' isAO34 tr1'NE tr2'NE unfolding tr1
  tr2 Opt.completedFrom-def by auto
    show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
    using ψ trn34 O33' O44' K1 ok unfolding ψ-def tr1 tr2
    using completedFromO-lastt less.premis(4,6)
    unfolding Opt.completedFrom-def using tr1 tr2 tr1'NE tr2'NE by auto
  next
  assume K2: ?K2
  then obtain sv1' where
  trn1: validTransV (sv1, sv1') and
  trn13: eqSec sv1 s1 and
  Atrn1: ¬ isIntV sv1 and ok': (statA' = statA ∨ statO = Diff) and
  Δ: Δ ∞ s1' s2' statA' sv1' sv2 statO
  unfolding match12-1-def by auto
  have r1': reachV sv1' using r1 trn1 by (metis Van.reach.Step fst-conv

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```

snd-conv)
  obtain trv1 trv2 where  $\psi: \psi s1' tr1' s2' tr2' statO sv1' trv1 sv2' trv2$ 
  using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r1' r2, simplified]
  using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
  using  $\psi O33' O44' tr1'NE tr2'NE$  unfolding tr1 tr2
  using Van.validFromS-Cons trn1 ok
  using K2 ok' Atrn1 eqSec-S-Cons trn13 trn34
  unfolding statA' Van.completedFrom-def eqSec-def
  using s13 tr1 tr1' tr2' trn3 trn4
by simp (smt (verit, ccfv-SIG) Opt.S.simps(2) Simple-Transition-System.lastt-Cons

Van.A.Cons-unfold Van.O.Cons-unfold  $\psi$ -def list.simps(3) status.simps(1))
next
  assume K3: ?K3
  then obtain sv2' where
    trn2: validTransV (sv2,sv2') and
    trn24: eqSec sv2 s2 and
    Atrn2:  $\neg isIntV sv2$  and  $ok': (statA' = statA \vee statO = Diff)$  and
     $\Delta: \Delta \infty s1' s2' statA' sv1 sv2' statO$ 
  unfolding match12-2-def by auto
  have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
  obtain trv1 trv2 where  $\psi: \psi s1' tr1' s2' tr2' statO sv1 trv1 sv2' trv2$ 
  using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r1 r2', simplified]
  using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
  show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\psi O33' O44' tr1'NE tr2'NE$  unfolding  $\psi$ -def tr1 tr2
  using Van.validFromS-Cons trn2 ok
  using K3 ok' Atrn2 eqSec-S-Cons trn24 trn34
  unfolding statA' Van.completedFrom-def
  by simp (metis last.simps lastt-def list.simps(3) status.distinct(2))
next
  assume K4: ?K4
  then obtain sv1' sv2' statO' where 0:
    statO' = sstatO' statO sv1 sv2
    validTransV (sv1,sv1')
    eqSec sv1 s1
    validTransV (sv2,sv2')
    eqSec sv2 s2
    Van.eqAct sv1 sv2
    and  $ok': statA' = Diff \longrightarrow statO' = Diff$  and  $\Delta: \Delta \infty s1' s2' statA'$ 
  sv1' sv2' statO'
  unfolding match12-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0
  by (metis Van.reach.Step fst-conv snd-conv)+
  obtain trv1 trv2 where  $\psi: \psi s1' tr1' s2' tr2' statO' sv1' trv1 sv2' trv2$ 

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```

    using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r12', simplified]
    using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.compleatedFrom-def by auto
    show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
    using trn34
    using  $\psi$  O33' O44' isAO34 tr1'NE tr2'NE unfolding  $\psi$ -def tr1 tr2
    using Van.validFromS-Cons 0
    using K4 eqSec-S-Cons
    unfolding statA' Van.eqAct-def Van.compleatedFrom-def match12-12-def
sstatO'-def
    by (smt (z3) Simple-Transition-System.lastt-Cons Van.A.Cons-unfold
Van.O.Cons-unfold
    compleatedFromO-lastt f3 f34 lastt-Nil less.premis(4) less.premis(6) list.inject
s13
    s24 status.simps(1) tr1 tr1' tr2 tr2' trn3 trn4 newStat.simps(1) new-
Stat.simps(3))
    qed
    qed
    qed
    qed
    qed

```

theorem *unwind-rsecure*:

```

assumes init: initCond  $\Delta$ 
and unwind: unwindCond  $\Delta$ 
shows rsecure
apply (rule rsecure-strong)
apply (elim conjE)
subgoal for s1 tr1 s2 tr2
    using init unfolding initCond-def
apply (erule-tac alle[of - s1])
apply (elim alle[of - s2] conjE)
apply (elim impE[of  $\langle$ istateO s1  $\wedge$  istateO s2 $\rangle$ ] exE conjE)
subgoal by clarify
subgoal for sv1 sv2
    using unwind apply (drule-tac unwindCond-ex- $\psi$ , blast+)
subgoal by (rule Transition-System.reach.Istate)
subgoal by (rule Transition-System.reach.Istate)
subgoal by (rule Transition-System.reach.Istate)
subgoal by (rule Transition-System.reach.Istate)
apply (elim exE)
subgoal for trv1 trv2
    apply (rule exI[of - sv1], rule exI[of - trv1], rule exI[of - sv2], rule exI[of -
trv2])
    by clarify
    .
    .
    .

```

3.4 Compositional unwinding

We allow networks of unwinding relations that unwind into each other, which offer a compositional alternative to monolithic unwinding.

definition *unwindIntoCond* ::

$$\begin{aligned} & (enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \Rightarrow \\ & (enat \Rightarrow 'stateO \Rightarrow 'stateO \Rightarrow status \Rightarrow 'stateV \Rightarrow 'stateV \Rightarrow status \Rightarrow bool) \\ & \Rightarrow bool \end{aligned}$$

where

$$\begin{aligned} & unwindIntoCond \Delta \Delta' \equiv \forall w s1 s2 statA sv1 sv2 statO. \\ & reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge \\ & \Delta w s1 s2 statA sv1 sv2 statO \longrightarrow \\ & (finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO \\ & s2) \\ & \wedge \\ & (statA = Eq \longrightarrow (isIntO s1 \longleftrightarrow isIntO s2)) \\ & \wedge \\ & ((\exists v < w. proact \Delta' v s1 s2 statA sv1 sv2 statO) \\ & \vee \\ & react \Delta' s1 s2 statA sv1 sv2 statO) \end{aligned}$$

lemma *unwindIntoCond-simpleI*:

assumes

$$\begin{aligned} & \bigwedge w s1 s2 statA sv1 sv2 statO. \\ & reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies \\ & \Delta w s1 s2 statA sv1 sv2 statO \\ & \implies \\ & (finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO \\ & s2) \end{aligned}$$

and

$$\begin{aligned} & \bigwedge w s1 s2 statA sv1 sv2 statO. \\ & reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies \\ & \Delta w s1 s2 statA sv1 sv2 statO \implies \\ & statA = Eq \end{aligned}$$

\implies

$$\begin{aligned} & isIntO s1 \longleftrightarrow isIntO s2 \\ & \bigwedge w s1 s2 statA sv1 sv2 statO. \\ & reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies \\ & \Delta w s1 s2 statA sv1 sv2 statO \end{aligned}$$

\implies

$$react \Delta' s1 s2 statA sv1 sv2 statO$$

shows *unwindIntoCond* $\Delta \Delta'$

using *assms unfolding unwindIntoCond-def by auto*

lemma *unwindIntoCond-simpleI2*:

assumes

$$\begin{aligned} & \bigwedge w s1 s2 statA sv1 sv2 statO. \\ & reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies \end{aligned}$$

$\Delta w s1 s2 statA sv1 sv2 statO$
 \implies
 $(finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO s2)$
and
 $\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO \implies$
 $statA = Eq$
 \implies
 $isIntO s1 \longleftrightarrow isIntO s2$
and
 $\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO$
 \implies
 $(\exists v < w. proact \Delta' v s1 s2 statA sv1 sv2 statO)$
shows *unwindIntoCond* $\Delta \Delta'$
using *assms unfolding unwindIntoCond-def by auto*

lemma *unwindIntoCond-simpleIB:*

assumes

$\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO$
 \implies
 $(finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO s2)$
and
 $\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO \implies$
 $statA = Eq$
 \implies
 $isIntO s1 \longleftrightarrow isIntO s2$
and
 $\bigwedge w s1 s2 statA sv1 sv2 statO.$
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$
 $\Delta w s1 s2 statA sv1 sv2 statO$
 \implies
 $(\exists v < w. proact \Delta' v s1 s2 statA sv1 sv2 statO) \vee react \Delta' s1 s2 statA sv1 sv2 statO$
shows *unwindIntoCond* $\Delta \Delta'$
using *assms unfolding unwindIntoCond-def by auto*

theorem *distrib-unwind-rsecure:*

assumes $m: 0 < m$ **and** $nxt: \bigwedge i. i < (m::nat) \implies nxt i \subseteq \{0..<m\}$

and *init:* *initCond* $(\Delta s 0)$

and *step:* $\bigwedge i. i < m \implies$

```

unwindIntoCond ( $\Delta s i$ ) ( $\lambda w s1 s2 statA sv1 sv2 statO$ .
   $\exists j \in \text{next } i. \Delta s j w s1 s2 statA sv1 sv2 statO$ )
shows rsecure
proof -
  define  $\Delta$  where  $D: \Delta \equiv \lambda w s1 s2 statA sv1 sv2 statO. \exists i < m. \Delta s i w s1 s2$ 
   $statA sv1 sv2 statO$ 
  have  $i: \text{initCond } \Delta$ 
  using  $\text{init } m$  unfolding  $\text{initCond-def } D$  by  $\text{meson}$ 
  have  $c: \text{unwindCond } \Delta$  unfolding  $\text{unwindCond-def}$  apply( $\text{intro allI impI allI}$ )
  apply( $\text{subst } (asm) D$ ) apply ( $\text{elim exE conjE}$ )
  subgoal for  $w s1 s2 statA sv1 sv2 statO i$ 
  apply( $\text{frule step}$ ) unfolding  $\text{unwindIntoCond-def}$ 
  apply( $\text{erule allE[of - w]}$ )
  apply( $\text{erule allE[of - s1]}$ ) apply( $\text{erule allE[of - s2]}$ ) apply( $\text{erule allE[of -$ 
   $statA]}$ )
  apply( $\text{erule allE[of - sv1]}$ ) apply( $\text{erule allE[of - sv2]}$ ) apply( $\text{erule allE[of -$ 
   $statO]}$ )
  apply  $\text{simp}$  apply( $\text{elim conjE}$ )
  apply( $\text{erule disjE}$ )
  subgoal apply( $\text{rule disjI1}$ )
  subgoal apply( $\text{elim exE conjE}$ ) subgoal for  $v$ 
  apply( $\text{rule exI[of - v], simp}$ )
  apply( $\text{rule proact-mono[of } \lambda w s1 s2 statA sv1 sv2 statO. \exists j \in \text{next } i. \Delta s j w$ 
   $s1 s2 statA sv1 sv2 statO]}$ )
  subgoal unfolding  $\text{le-fun-def } D$  by  $\text{simp}$  ( $\text{meson atLeastLessThan-iff next}$ 
   $\text{subsetD}$ )
  subgoal . . . .
  subgoal apply( $\text{rule disjI2}$ )
  apply( $\text{rule match-mono[of } \lambda w s1 s2 statA sv1 sv2 statO. \exists j \in \text{next } i. \Delta s j w$ 
   $s1 s2 statA sv1 sv2 statO]}$ )
  subgoal unfolding  $\text{le-fun-def } D$  by  $\text{simp}$  ( $\text{meson atLeastLessThan-iff next}$ 
   $\text{subsetD}$ )
  subgoal . . . .
  show  $?thesis$  using  $\text{unwind-rsecure[OF } i c]$  .
qed

corollary  $\text{linear-unwind-rsecure}$ :
assumes  $\text{init: initCond } (\Delta s 0)$ 
and  $\text{step: } (\bigwedge i. i < m \implies$ 
   $\text{unwindIntoCond } (\Delta s i) (\lambda w s1 s2 statA sv1 sv2 statO$ .
     $\Delta s i w s1 s2 statA sv1 sv2 statO \vee$ 
     $\Delta s (\text{Suc } i) w s1 s2 statA sv1 sv2 statO))$ 
and  $\text{finish: unwindIntoCond } (\Delta s m) (\Delta s m)$ 
shows rsecure
apply( $\text{rule distrib-unwind-rsecure[of Suc } m \lambda i. \text{if } i < m \text{ then } \{i, \text{Suc } i\} \text{ else } \{i\} \Delta s,$ 
   $\text{OF - - init}]$ )
using  $\text{step finish}$ 
by ( $\text{auto simp: less-Suc-eq-le}$ )

```

definition *oor* where

$oor \Delta \Delta_2 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$

$\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO$

lemma *oorI1*:

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor \Delta \Delta_2 w s1 s2 statA sv1 sv2 statO$

unfolding *oor-def* by *simp*

lemma *oorI2*:

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor \Delta \Delta_2 w s1 s2 statA sv1 sv2 statO$

unfolding *oor-def* by *simp*

definition *oor3* where

$oor3 \Delta \Delta_2 \Delta_3 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$

$\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$

$\Delta_3 w s1 s2 statA sv1 sv2 statO$

lemma *oor3I1*:

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w s1 s2 statA sv1 sv2 statO$

unfolding *oor3-def* by *simp*

lemma *oor3I2*:

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w s1 s2 statA sv1 sv2 statO$

unfolding *oor3-def* by *simp*

lemma *oor3I3*:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w s1 s2 statA sv1 sv2 statO$

unfolding *oor3-def* by *simp*

definition *oor4* where

$oor4 \Delta \Delta_2 \Delta_3 \Delta_4 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$

$\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$

$\Delta_3 w s1 s2 statA sv1 sv2 statO \vee \Delta_4 w s1 s2 statA sv1 sv2 statO$

lemma *oor4I1*:

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I2*:

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I3*:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I4*:

$\Delta_4 w s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w s1 s2 statA sv1 sv2 statO$
unfolding *oor4-def by simp*

definition *oor5 where*

$oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w s1 s2 statA sv1 sv2 statO \vee \Delta_4 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_5 w s1 s2 statA sv1 sv2 statO$

lemma *oor5I1:*

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def by simp*

lemma *oor5I2:*

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def by simp*

lemma *oor5I3:*

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def by simp*

lemma *oor5I4:*

$\Delta_4 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def by simp*

lemma *oor5I5:*

$\Delta_5 w s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def by simp*

definition *oor6 where*

$oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 \equiv \lambda w s1 s2 statA sv1 sv2 statO.$
 $\Delta w s1 s2 statA sv1 sv2 statO \vee \Delta_2 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w s1 s2 statA sv1 sv2 statO \vee \Delta_4 w s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_5 w s1 s2 statA sv1 sv2 statO \vee \Delta_6 w s1 s2 statA sv1 sv2 statO$

lemma *oor6I1:*

$\Delta w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def by simp*

lemma *oor6I2:*

$\Delta_2 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def by simp*

lemma *oor6I3*:

$\Delta_3 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* **by** *simp*

lemma *oor6I4*:

$\Delta_4 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* **by** *simp*

lemma *oor6I5*:

$\Delta_5 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* **by** *simp*

lemma *oor6I6*:

$\Delta_6 w s1 s2 statA sv1 sv2 statO \implies oor6 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \Delta_6 w s1 s2 statA sv1 sv2 statO$

unfolding *oor6-def* **by** *simp*

lemma *unwind-rsecure-foo*:

assumes *init*: *initCond* Δ_0

and *step0*: *unwindIntoCond* $\Delta_0 \Delta NN$

and *stepNN*: *unwindIntoCond* ΔNN (*oor5* $\Delta NN \Delta SN \Delta NS \Delta SS \Delta nonspec$)

and *stepNS*: *unwindIntoCond* ΔNS (*oor4* $\Delta NN \Delta SN \Delta NS \Delta SS$)

and *stepSN*: *unwindIntoCond* ΔSN (*oor4* $\Delta NN \Delta SN \Delta NS \Delta SS$)

and *stepSS*: *unwindIntoCond* ΔSS (*oor4* $\Delta NN \Delta SN \Delta NS \Delta SS$)

and *stepNonspec*: *unwindIntoCond* $\Delta nonspec \Delta nonspec$

shows *rsecure*

proof–

define *m* **where** *m*: $m \equiv (6::nat)$

define Δs **where** Δs : $\Delta s \equiv \lambda i::nat.$

if $i = 0$ *then* Δ_0

else if $i = 1$ *then* ΔNN

else if $i = 2$ *then* ΔSN

else if $i = 3$ *then* ΔNS

else if $i = 4$ *then* ΔSS

else $\Delta nonspec$

define *next* **where** *next*: $next \equiv \lambda i::nat.$

if $i = 0$ *then* $\{1::nat\}$

else if $i = 1$ *then* $\{1,2,3,4,5\}$

else if $i = 2$ *then* $\{1,2,3,4\}$

else if $i = 3$ *then* $\{1,2,3,4\}$

else if $i = 4$ *then* $\{1,2,3,4\}$

else $\{5\}$

```

show ?thesis apply(rule distrib-unwind-rsecure[of m nat  $\Delta$ s])
  subgoal unfolding m by auto
  subgoal unfolding nat m by auto
  subgoal using init unfolding  $\Delta$ s by auto
  subgoal
    unfolding m nat  $\Delta$ s
    using step0 stepNN stepNS stepSN stepSS stepNonspec
    unfolding oor4-def oor5-def by auto .
qed

```

```

lemma isIntO-match1: isIntO s1  $\implies$  match1  $\Delta$  s1 s2 statA sv1 sv2 statO
unfolding match1-def by auto

```

```

lemma isIntO-match2: isIntO s2  $\implies$  match2  $\Delta$  s1 s2 statA sv1 sv2 statO
unfolding match2-def by auto

```

```

lemma match1-1-oorI1:
  match1-1  $\Delta$  s1 s1' s2 statA sv1 sv2 statO  $\implies$ 
  match1-1 (oor  $\Delta$   $\Delta_2$ ) s1 s1' s2 statA sv1 sv2 statO
apply(rule match1-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match1-1-oorI2:
  match1-1  $\Delta_2$  s1 s1' s2 statA sv1 sv2 statO  $\implies$ 
  match1-1 (oor  $\Delta$   $\Delta_2$ ) s1 s1' s2 statA sv1 sv2 statO
apply(rule match1-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match1-oorI1:
  match1  $\Delta$  s1 s2 statA sv1 sv2 statO  $\implies$ 
  match1 (oor  $\Delta$   $\Delta_2$ ) s1 s2 statA sv1 sv2 statO
apply(rule match1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match1-oorI2:
  match1  $\Delta_2$  s1 s2 statA sv1 sv2 statO  $\implies$ 
  match1 (oor  $\Delta$   $\Delta_2$ ) s1 s2 statA sv1 sv2 statO
apply(rule match1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match2-1-oorI1:
  match2-1  $\Delta$  s1 s2 s2' statA sv1 sv2 statO  $\implies$ 
  match2-1 (oor  $\Delta$   $\Delta_2$ ) s1 s2 s2' statA sv1 sv2 statO
apply(rule match2-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match2-1-oorI2:
  match2-1  $\Delta_2$  s1 s2 s2' statA sv1 sv2 statO  $\implies$ 

```

match2-1 (oor $\Delta \Delta_2$) s1 s2 s2' statA sv1 sv2 statO
apply(rule *match2-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-oorI1*:
match2 Δ s1 s2 statA sv1 sv2 statO \implies
match2 (oor $\Delta \Delta_2$) s1 s2 statA sv1 sv2 statO
apply(rule *match2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match2-oorI2*:
match2 Δ_2 s1 s2 statA sv1 sv2 statO \implies
match2 (oor $\Delta \Delta_2$) s1 s2 statA sv1 sv2 statO
apply(rule *match2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-oorI1*:
match12 Δ s1 s2 statA sv1 sv2 statO \implies
match12 (oor $\Delta \Delta_2$) s1 s2 statA sv1 sv2 statO
apply(rule *match12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-oorI2*:
match12 Δ_2 s1 s2 statA sv1 sv2 statO \implies
match12 (oor $\Delta \Delta_2$) s1 s2 statA sv1 sv2 statO
apply(rule *match12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-1-oorI1*:
match12-1 Δ s1' s2' statA' sv1 sv2 statO \implies
match12-1 (oor $\Delta \Delta_2$) s1' s2' statA' sv1 sv2 statO
apply(rule *match12-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-1-oorI2*:
match12-1 Δ_2 s1' s2' statA' sv1 sv2 statO \implies
match12-1 (oor $\Delta \Delta_2$) s1' s2' statA' sv1 sv2 statO
apply(rule *match12-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-2-oorI1*:
match12-2 Δ s2 s2' statA' sv1 sv2 statO \implies
match12-2 (oor $\Delta \Delta_2$) s2 s2' statA' sv1 sv2 statO
apply(rule *match12-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-2-oorI2*:
match12-2 Δ_2 s2 s2' statA' sv1 sv2 statO \implies
match12-2 (oor $\Delta \Delta_2$) s2 s2' statA' sv1 sv2 statO
apply(rule *match12-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-12-oorI1*:
match12-12 Δ s1' s2' statA' sv1 sv2 statO \implies
match12-12 (oor $\Delta \Delta_2$) s1' s2' statA' sv1 sv2 statO
apply(rule *match12-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match12-12-oorI2*:
 $match12-12 \Delta_2 s1' s2' statA' sv1 sv2 statO \implies$
 $match12-12 (oor \Delta \Delta_2) s1' s2' statA' sv1 sv2 statO$
apply(rule *match12-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI1*:
 $react \Delta s1 s2 statA sv1 sv2 statO \implies$
 $react (oor \Delta \Delta_2) s1 s2 statA sv1 sv2 statO$
apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI2*:
 $react \Delta_2 s1 s2 statA sv1 sv2 statO \implies$
 $react (oor \Delta \Delta_2) s1 s2 statA sv1 sv2 statO$
apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI1*:
 $proact \Delta meas s1 s2 statA sv1 sv2 statO \implies$
 $proact (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$
apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI2*:
 $proact \Delta_2 meas s1 s2 statA sv1 sv2 statO \implies$
 $proact (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$
apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-1-oorI1*:
 $move-1 \Delta meas s1 s2 statA sv1 sv2 statO \implies$
 $move-1 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$
apply(rule *move-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-1-oorI2*:
 $move-1 \Delta_2 meas s1 s2 statA sv1 sv2 statO \implies$
 $move-1 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$
apply(rule *move-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-2-oorI1*:
 $move-2 \Delta meas s1 s2 statA sv1 sv2 statO \implies$
 $move-2 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$
apply(rule *move-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-2-oorI2*:
 $move-2 \Delta_2 meas s1 s2 statA sv1 sv2 statO \implies$
 $move-2 (oor \Delta \Delta_2) meas s1 s2 statA sv1 sv2 statO$
apply(rule *move-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-12-oorI1*:
 $move-12 \Delta meas\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$
 $move-12\ (oor\ \Delta\ \Delta_2)\ meas\ s1\ s2\ statA\ sv1\ sv2\ statO$
apply(rule *move-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-12-oorI2*:
 $move-12\ \Delta_2\ meas\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$
 $move-12\ (oor\ \Delta\ \Delta_2)\ meas\ s1\ s2\ statA\ sv1\ sv2\ statO$
apply(rule *move-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

end

context *Relative-Security-Determ*
begin

lemma *match1-1-defD*: $match1-1\ \Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \longleftrightarrow$
 $\neg\ finalV\ sv1 \wedge \Delta \infty\ s1'\ s2\ statA\ (nextO\ sv1)\ sv2\ statO$
unfolding *match1-1-def validTrans-iff-next* **by** *simp*

lemma *match1-12-defD*: $match1-12\ \Delta\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \longleftrightarrow$
 $\neg\ finalV\ sv1 \wedge \neg\ finalV\ sv2 \wedge$
 $\Delta \infty\ s1'\ s2\ statA\ (nextO\ sv1)\ (nextO\ sv2)\ (sstatO'\ statO\ sv1\ sv2)$
unfolding *match1-12-def validTrans-iff-next* **by** *simp*

lemmas *match1-defsD* = *match1-def match1-1-defD match1-12-defD*

lemma *match2-1-defD*: $match2-1\ \Delta\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO \longleftrightarrow$
 $\neg\ finalV\ sv2 \wedge \Delta \infty\ s1\ s2'\ statA\ sv1\ (nextO\ sv2)\ statO$
unfolding *match2-1-def validTrans-iff-next* **by** *simp*

lemma *match2-12-defD*: $match2-12\ \Delta\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO \longleftrightarrow$
 $\neg\ finalV\ sv1 \wedge \neg\ finalV\ sv2 \wedge \Delta \infty\ s1\ s2'\ statA\ (nextO\ sv1)\ (nextO\ sv2)\ (sstatO'$
 $statO\ sv1\ sv2)$
unfolding *match2-12-def validTrans-iff-next* **by** *simp*

lemmas *match2-defsD* = *match2-def match2-1-defD match2-12-defD*

lemma *match12-1-defD*: $match12-1\ \Delta\ s1'\ s2'\ statA'\ sv1\ sv2\ statO \longleftrightarrow$
 $\neg\ finalV\ sv1 \wedge \Delta \infty\ s1'\ s2'\ statA'\ (nextO\ sv1)\ sv2\ statO$
unfolding *match12-1-def validTrans-iff-next* **by** *simp*

lemma *match12-2-defD*: $match12-2\ \Delta\ s1'\ s2'\ statA'\ sv1\ sv2\ statO \longleftrightarrow$
 $\neg\ finalV\ sv2 \wedge \Delta \infty\ s1'\ s2'\ statA'\ sv1\ (nextO\ sv2)\ statO$

unfolding *match12-2-def validTrans-iff-next by simp*

lemma *match12-12-defD*: $match12-12 \Delta s1' s2' statA' sv1 sv2 statO \longleftrightarrow$
(*let* $statO' = sstatO' statO sv1 sv2$ *in*
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $(statA' = Diff \longrightarrow statO' = Diff) \wedge$
 $\Delta \infty s1' s2' statA' (nextO sv1) (nextO sv2) statO'$)

unfolding *match12-12-def validTrans-iff-next by simp*

lemmas *match12-defsD = match12-def match12-1-defD match12-2-defD match12-12-defD*

lemmas *match-deep-defsD = match1-defsD match2-defsD match12-defsD*

lemma *move-1-defD*: $move-1 \Delta w s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv1 \wedge \Delta w s1 s2 statA (nextO sv1) sv2 statO$

unfolding *move-1-def validTrans-iff-next by simp*

lemma *move-2-defD*: $move-2 \Delta w s1 s2 statA sv1 sv2 statO \longleftrightarrow$
 $\neg finalV sv2 \wedge \Delta w s1 s2 statA sv1 (nextO sv2) statO$

unfolding *move-2-def validTrans-iff-next by simp*

lemma *move-12-defD*: $move-12 \Delta w s1 s2 statA sv1 sv2 statO \longleftrightarrow$
(*let* $statO' = sstatO' statO sv1 sv2$ *in*
 $\neg finalV sv1 \wedge \neg finalV sv2 \wedge$
 $\Delta w s1 s2 statA (nextO sv1) (nextO sv2) statO'$)

unfolding *move-12-def validTrans-iff-next by simp*

lemmas *proact-defsD = proact-def move-1-defD move-2-defD move-12-defD*

end

end

4 Unwinding Proof Method for Relative Security

This theory formalizes the notion of unwinding for relative security, and proves its soundness.

theory *Unwinding*
imports *Relative-Security*
begin

4.1 The types and operators underlying unwinding: status, matching operators, etc.

context *Rel-Sec*

begin

datatype *status* = *Eq* | *Diff*

fun *newStat* :: *status* \Rightarrow *bool* \times '*a* \Rightarrow *bool* \times '*a* \Rightarrow *status* **where**
 newStat Eq (*True*,*a*) (*True*,*a'*) = (*if a = a' then Eq else Diff*)
 newStat stat - - = *stat*

definition *sstatO'* *statO* *sv1* *sv2* = *newStat statO* (*isIntV sv1*, *getObsV sv1*)
(*isIntV sv2*, *getObsV sv2*)

definition *sstatA'* *statA* *s1* *s2* = *newStat statA* (*isIntO s1*, *getObsO s1*) (*isIntO*
s2, *getObsO s2*)

lemma *newStat-EqI*:

assumes $\langle R = S \rangle$

shows $\langle \text{newStat Eq } (P, R) (Q, S) = \text{Eq} \rangle$

apply (*cases P*)

apply (*metis assms newStat.simps(1) newStat.simps(4)*)

by (*cases Q*) *auto*

lemma *newStat-diff*: *newStat stat r r = Diff* \Longrightarrow *stat = Diff*

by (*metis newStat.elims newStat.simps(1)*)

definition *initCond* ::

(*enat* \Rightarrow *enat* \Rightarrow *enat* \Rightarrow '*stateO* \Rightarrow '*stateO* \Rightarrow *status* \Rightarrow '*stateV* \Rightarrow '*stateV* \Rightarrow
status \Rightarrow *bool*) \Rightarrow *bool* **where**

initCond $\Delta \equiv \forall s1\ s2.$

istateO s1 \wedge *istateO s2*

\longrightarrow

($\exists sv1\ sv2. \text{istateV } sv1 \wedge \text{istateV } sv2 \wedge \text{corrState } sv1\ s1 \wedge \text{corrState } sv2\ s2$
 $\wedge \Delta \infty \infty \infty s1\ s2 \text{Eq } sv1\ sv2 \text{Eq}$)

definition *match1-1* $\Delta\ w1\ w2\ s1\ s1'\ s2\ \text{statA}\ sv1\ sv2\ \text{statO} \equiv$

$\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$

$\Delta \infty w1\ w2\ s1'\ s2\ \text{statA}\ sv1'\ sv2\ \text{statO}$

definition *match1-12* $\Delta\ w1\ w2\ s1\ s1'\ s2\ \text{statA}\ sv1\ sv2\ \text{statO} \equiv$

($\exists sv1'\ sv2'.$

let *statO'* = *sstatO' statO sv1 sv2* *in*

validTransV (*sv1*,*sv1'*) \wedge

validTransV (*sv2*,*sv2'*) \wedge

$\Delta \infty w1\ w2\ s1'\ s2\ \text{statA}\ sv1'\ sv2'\ \text{statO}'$)

definition $match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $\neg isIntO s1 \longrightarrow$
 $(\forall s1'. validTransO (s1, s1'))$
 \longrightarrow
 $(\exists w1' < w1. \exists w2' < w2. \neg isSecO s1 \wedge \Delta \infty w1' w2' s1' s2 statA sv1 sv2$
 $statO) \vee$
 $(\exists w2' < w2. eqSec sv1 s1 \wedge \neg isIntV sv1 \wedge match1-1 \Delta \infty w2' s1 s1' s2$
 $statA sv1 sv2 statO) \vee$
 $(eqSec sv1 s1 \wedge \neg isSecV sv2 \wedge Van.eqAct sv1 sv2 \wedge match1-12 \Delta \infty \infty s1$
 $s1' s2 statA sv1 sv2 statO))$

lemmas $match1-defs = match1-def match1-1-def match1-12-def$

lemma $match1-1-mono$:

$\Delta \leq \Delta' \implies match1-1 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \implies$
 $match1-1 \Delta' w1 w2 s1 s1' s2 statA sv1 sv2 statO$

unfolding $le-fun-def match1-1-def$ **by** $auto$

lemma $match1-12-mono$:

$\Delta \leq \Delta' \implies match1-12 \Delta w1 w2 s1 s1' s2 statA sv1 sv2 statO \implies$
 $match1-12 \Delta' w1 w2 s1 s1' s2 statA sv1 sv2 statO$

unfolding $le-fun-def match1-12-def$ **by** $fastforce$

lemma $match1-mono$:

assumes $\Delta \leq \Delta'$

shows $match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies match1 \Delta' w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding $match1-def$ **apply** $clarify$ **subgoal for** $s1'$ **apply** $(erule alle[of - s1 \uparrow])$

using $match1-1-mono[OF assms, of - - s1 s1' s2 statA sv1 sv2 statO]$

$match1-12-mono[OF assms, of - - s1 s1' s2 statA sv1 sv2 statO]$

$assms[unfolding le-fun-def, rule-format, of - - - s1' s2 statA sv1 sv2 statO]$

by $fastforce$.

definition $match2-1 \Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \equiv$

$\exists sv2'. validTransV (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1 s2' statA sv1 sv2' statO$

definition $match2-12 \Delta w1 w2 s1 s2 s2' statA sv1 sv2 statO \equiv$

$\exists sv1' sv2'.$
 $let statO' = sstatO' statO sv1 sv2 in$
 $validTransV (sv1, sv1') \wedge$
 $validTransV (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1 s2' statA sv1' sv2' statO'$

definition $match2 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$

$\neg isIntO s2 \longrightarrow$

$(\forall s2'. \text{validTransO } (s2, s2'))$
 \longrightarrow
 $(\exists w1' < w1. \exists w2' < w2. \neg \text{isSecO } s2 \wedge \Delta \infty w1' w2' s1 s2' \text{statA } sv1 sv2 \text{statO}) \vee$
 $(\exists w1' < w1. \text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge \text{match2-1 } \Delta w1' \infty s1 s2 s2' \text{statA } sv1 sv2 \text{statO}) \vee$
 $(\neg \text{isSecV } sv1 \wedge \text{eqSec } sv2 s2 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match2-12 } \Delta \infty \infty s1 s2 s2' \text{statA } sv1 sv2 \text{statO}))$

lemmas *match2-defs* = *match2-def match2-1-def match2-12-def*

lemma *match2-1-mono*:

$\Delta < \Delta' \implies \text{match2-1 } \Delta w1 w2 s1 s1' s2 \text{statA } sv1 sv2 \text{statO} \implies \text{match2-1 } \Delta' w1 w2 s1 s1' s2 \text{statA } sv1 sv2 \text{statO}$

unfolding *le-fun-def match2-1-def* **by** *auto*

lemma *match2-12-mono*:

$\Delta \leq \Delta' \implies \text{match2-12 } \Delta w1 w2 s1 s1' s2 \text{statA } sv1 sv2 \text{statO} \implies \text{match2-12 } \Delta' w1 w2 s1 s1' s2 \text{statA } sv1 sv2 \text{statO}$

unfolding *le-fun-def match2-12-def* **by** *fastforce*

lemma *match2-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{match2 } \Delta w1 w2 s1 s2 \text{statA } sv1 sv2 \text{statO} \implies \text{match2 } \Delta' w1 w2 s1 s2 \text{statA } sv1 sv2 \text{statO}$

unfolding *match2-def* **apply** *clarify subgoal for s2'* **apply**(*erule alle[of - s2']*)

using *match2-1-mono[OF assms, of - - s1 s2 s2' statA sv1 sv2 statO]*

match2-12-mono[OF assms, of - - s1 s2 s2' statA sv1 sv2 statO]

assms[unfolded le-fun-def, rule-format, of - - - s1 s2' statA sv1 sv2 statO]

by *fastforce* .

definition *match12-1* $\Delta w1 w2 s1' s2' \text{statA}' sv1 sv2 \text{statO} \equiv$

$\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$
 $\Delta \infty w1 w2 s1' s2' \text{statA}' sv1' sv2 \text{statO}$

definition *match12-2* $\Delta w1 w2 s1' s2' \text{statA}' sv1 sv2 \text{statO} \equiv$

$\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$
 $\Delta \infty w1 w2 s1' s2' \text{statA}' sv1 sv2' \text{statO}$

definition *match12-12* $\Delta w1 w2 s1' s2' \text{statA}' sv1 sv2 \text{statO} \equiv$

$\exists sv1' sv2'.$
let $\text{statO}' = \text{sstatO}' \text{statO } sv1 sv2$ *in*
 $\text{validTransV } (sv1, sv1') \wedge$
 $\text{validTransV } (sv2, sv2') \wedge$
 $(\text{statA}' = \text{Diff} \longrightarrow \text{statO}' = \text{Diff}) \wedge$
 $\Delta \infty w1 w2 s1' s2' \text{statA}' sv1' sv2' \text{statO}'$

definition $match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $\forall s1' s2'.$
 $let\ statA' = sstatA' statA\ s1\ s2\ in$
 $validTransO\ (s1, s1') \wedge$
 $validTransO\ (s2, s2') \wedge$
 $Opt.eqAct\ s1\ s2 \wedge$
 $isIntO\ s1 \wedge isIntO\ s2$
 \longrightarrow
 $(\exists w1' < w1. \exists w2' < w2. \neg isSecO\ s1 \wedge \neg isSecO\ s2 \wedge (statA = statA' \vee statO = Diff)) \wedge$
 $\Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\exists w2' < w2. \neg isSecO\ s2 \wedge eqSec\ sv1\ s1 \wedge \neg isIntV\ sv1 \wedge$
 $(statA = statA' \vee statO = Diff) \wedge$
 $match12-1\ \Delta \infty w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(\exists w1' < w1. \neg isSecO\ s1 \wedge eqSec\ sv2\ s2 \wedge \neg isIntV\ sv2 \wedge$
 $(statA = statA' \vee statO = Diff) \wedge$
 $match12-2\ \Delta w1' \infty s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(eqSec\ sv1\ s1 \wedge eqSec\ sv2\ s2 \wedge Van.eqAct\ sv1\ sv2 \wedge$
 $match12-12\ \Delta \infty \infty s1' s2' statA' sv1 sv2 statO)$

lemmas $match12-defs = match12-def\ match12-1-def\ match12-2-def\ match12-12-def$

lemma $match12-simpleI:$
assumes $\bigwedge s1' s2' statA'.$
 $statA' = sstatA' statA\ s1\ s2 \implies$
 $validTransO\ (s1, s1') \implies$
 $validTransO\ (s2, s2') \implies$
 $Opt.eqAct\ s1\ s2 \implies$
 $(\exists w1' < w1. \exists w2' < w2. \neg isSecO\ s1 \wedge \neg isSecO\ s2 \wedge (statA = statA' \vee statO = Diff)) \wedge$
 $\Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO)$
 \vee
 $(eqSec\ sv1\ s1 \wedge eqSec\ sv2\ s2 \wedge Van.eqAct\ sv1\ sv2 \wedge$
 $match12-12\ \Delta \infty \infty s1' s2' statA' sv1 sv2 statO)$
shows $match12 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$
using $assms\ unfolding\ match12-def\ Let-def\ by\ blast$

lemma $match12-1-mono:$
 $\Delta \leq \Delta' \implies match12-1\ \Delta\ w1\ w2\ s1'\ s2'\ statA'\ sv1\ sv2\ statO \implies match12-1\ \Delta'\ w1\ w2\ s1'\ s2'\ statA'\ sv1\ sv2\ statO$
unfolding $le-fun-def\ match12-1-def\ by\ auto$

lemma $match12-2-mono:$
 $\Delta \leq \Delta' \implies match12-2\ \Delta\ w1\ w2\ s1\ s2'\ statA'\ sv1\ sv2\ statO \implies match12-2\ \Delta'\ w1\ w2\ s1\ s2'\ statA'\ sv1\ sv2\ statO$

unfolding *le-fun-def match12-2-def* by *auto*

lemma *match12-12-mono*:

$\Delta \leq \Delta' \implies \text{match12-12 } \Delta \ w1 \ w2 \ s1' \ s2' \ \text{statA}' \ sv1 \ sv2 \ \text{statO} \implies \text{match12-12}$
 $\Delta' \ w1 \ w2 \ s1' \ s2' \ \text{statA}' \ sv1 \ sv2 \ \text{statO}$

unfolding *le-fun-def match12-12-def* by *fastforce*

lemma *match12-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{match12 } \Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \implies \text{match12 } \Delta' \ w1 \ w2 \ s1 \ s2$
 $\text{statA} \ sv1 \ sv2 \ \text{statO}$

unfolding *match12-def* **apply** *clarify* **subgoal** **for** $s1' \ s2'$ **apply**(*erule* *allE*[*of* -
 $s1'$]) **apply**(*erule* *allE*[*of* - $s2'$])

using *match12-1-mono*[*OF* *assms*, *of* - - $s1' \ s2' - sv1 \ sv2 \ \text{statO}$]

match12-2-mono[*OF* *assms*, *of* - - $s1' \ s2' - sv1 \ sv2 \ \text{statO}$]

match12-12-mono[*OF* *assms*, *of* - - $s1' \ s2' - sv1 \ sv2 \ \text{statO}$]

assms[*unfolded* *le-fun-def*, *rule-format*, *of* - - - $s1' \ s2'$
 $s\text{statA}' \ \text{statA} \ s1 \ s2 \ sv1 \ sv2 \ \text{statO}$]

apply *simp* by *blast* .

definition *react* $\Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \equiv$

$\text{match1 } \Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$

\wedge

$\text{match2 } \Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$

\wedge

$\text{match12 } \Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$

lemmas *react-defs* = *match1-def* *match2-def* *match12-def*

lemmas *match-deep-defs* = *match1-defs* *match2-defs* *match12-defs*

lemma *match-mono*:

assumes $\Delta \leq \Delta'$

shows $\text{react } \Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \implies \text{react } \Delta' \ w1 \ w2 \ s1 \ s2 \ \text{statA}$
 $sv1 \ sv2 \ \text{statO}$

unfolding *react-def* **using** *match1-mono*[*OF* *assms*] *match2-mono*[*OF* *assms*] *match12-mono*[*OF*
assms] by *auto*

definition *move-1* $\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \equiv$

$\exists sv1'. \text{validTransV } (sv1, sv1') \wedge$

$\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1' \ sv2 \ \text{statO}$

definition *move-2* $\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \equiv$

$\exists sv2'. \text{validTransV } (sv2, sv2') \wedge$

$\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2' \ \text{statO}$

definition *move-12* $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $\exists sv1' sv2'$.
let $statO' = sstatO' statO sv1 sv2$ *in*
 $validTransV (sv1, sv1') \wedge validTransV (sv2, sv2') \wedge$
 $\Delta w w1 w2 s1 s2 statA sv1' sv2' statO'$

definition *proact* $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \equiv$
 $(\neg isSecV sv1 \wedge \neg isIntV sv1 \wedge move-1 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO)$
 \vee
 $(\neg isSecV sv2 \wedge \neg isIntV sv2 \wedge move-2 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO)$
 \vee
 $(\neg isSecV sv1 \wedge \neg isSecV sv2 \wedge Van.eqAct sv1 sv2 \wedge move-12 \Delta w w1 w2 s1 s2$
 $statA sv1 sv2 statO)$

lemmas *proact-defs* = *proact-def move-1-def move-2-def move-12-def*

lemma *move-1-mono*:

$\Delta \leq \Delta' \implies move-1 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies move-1 \Delta' w w1$
 $w2 s1 s2 statA sv1 sv2 statO$

unfolding *le-fun-def move-1-def* **by** *auto*

lemma *move-2-mono*:

$\Delta \leq \Delta' \implies move-2 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies move-2 \Delta' w w1$
 $w2 s1 s2 statA sv1 sv2 statO$

unfolding *le-fun-def move-2-def* **by** *auto*

lemma *move-12-mono*:

$\Delta \leq \Delta' \implies move-12 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies move-12 \Delta' w w1$
 $w2 s1 s2 statA sv1 sv2 statO$

unfolding *le-fun-def move-12-def* **by** *fastforce*

lemma *proact-mono*:

assumes $\Delta \leq \Delta'$

shows *proact* $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies proact \Delta' w w1 w2 s1 s2$
 $statA sv1 sv2 statO$

unfolding *proact-def* **using** *move-1-mono[OF assms] move-2-mono[OF assms]*
move-12-mono[OF assms] **by** *auto*

4.2 The definition of unwinding

definition *unwindCond* ::

$(enat \implies enat \implies enat \implies 'stateO \implies 'stateO \implies status \implies 'stateV \implies 'stateV \implies$
 $status \implies bool) \implies bool$

where

unwindCond $\Delta \equiv \forall w w1 w2 s1 s2 statA sv1 sv2 statO.$

$reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$

\longrightarrow

$(finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO$

$$\begin{aligned}
& s2) \\
& \wedge \\
& (statA = Eq \longrightarrow (isIntO\ s1 \longleftrightarrow isIntO\ s2)) \\
& \wedge \\
& ((\exists v < w. \text{proact } \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO) \\
& \vee \\
& \text{react } \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \\
&)
\end{aligned}$$

lemma *unwindCond-simpleI*:

assumes

$$\begin{aligned}
& \wedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO. \\
& \text{reachO } s1 \Longrightarrow \text{reachO } s2 \Longrightarrow \text{reachV } sv1 \Longrightarrow \text{reachV } sv2 \Longrightarrow \\
& \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \\
& \Longrightarrow \\
& (\text{finalO } s1 \longleftrightarrow \text{finalO } s2) \wedge (\text{finalV } sv1 \longleftrightarrow \text{finalO } s1) \wedge (\text{finalV } sv2 \longleftrightarrow \text{finalO } \\
& s2)
\end{aligned}$$

and

$$\begin{aligned}
& \wedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO. \\
& \text{reachO } s1 \Longrightarrow \text{reachO } s2 \Longrightarrow \text{reachV } sv1 \Longrightarrow \text{reachV } sv2 \Longrightarrow \\
& \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \Longrightarrow statA = Eq \\
& \Longrightarrow \\
& isIntO\ s1 \longleftrightarrow isIntO\ s2
\end{aligned}$$

and

$$\begin{aligned}
& \wedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO. \\
& \text{reachO } s1 \Longrightarrow \text{reachO } s2 \Longrightarrow \text{reachV } sv1 \Longrightarrow \text{reachV } sv2 \Longrightarrow \\
& \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \\
& \Longrightarrow \\
& \text{react } \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO
\end{aligned}$$

shows *unwindCond* Δ

using *assms* **unfolding** *unwindCond-def* **by** *auto*

4.3 The soundness of unwinding

The proof of soundness for general unwinding is significantly more elaborate than that for the finitary case.

definition $\psi\ s1\ tr1\ s2\ tr2\ statO\ sv1\ trv1\ sv2\ trv2 \equiv$

$$\begin{aligned}
& trv1 \neq [] \wedge trv2 \neq [] \wedge \\
& Van.\text{validFromS } sv1\ trv1 \wedge \\
& Van.\text{validFromS } sv2\ trv2 \wedge \\
& (\text{finalV } (\text{lastt } sv1\ trv1) \longleftrightarrow \text{finalO } (\text{lastt } s1\ tr1)) \wedge (\text{finalV } (\text{lastt } sv2\ trv2) \longleftrightarrow \\
& \text{finalO } (\text{lastt } s2\ tr2)) \wedge \\
& Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge \\
& Van.A\ trv1 = Van.A\ trv2 \wedge \\
& (statO = Eq \wedge Opt.O\ tr1 \neq Opt.O\ tr2 \longrightarrow Van.O\ trv1 \neq Van.O\ trv2)
\end{aligned}$$

lemma ψ -completedFrom: completedFromO s1 tr1 \implies completedFromO s2 tr2 \implies

ψ s1 tr1 s2 tr2 statO sv1 trv1 sv2 trv2
 \implies completedFromV sv1 trv1 \wedge completedFromV sv2 trv2

unfolding ψ -def Opt.completedFrom-def Van.completedFrom-def lastt-def
by presburger

lemma completedFromO-lastt: completedFromO s1 tr1 \implies finalO (lastt s1 tr1)

unfolding Opt.completedFrom-def lastt-def **by** auto

lemma rsecure-strong:

assumes

\wedge s1 tr1 s2 tr2.

istateO s1 \wedge Opt.validFromS s1 tr1 \wedge completedFromO s1 tr1 \wedge

istateO s2 \wedge Opt.validFromS s2 tr2 \wedge completedFromO s2 tr2 \wedge

Opt.A tr1 = Opt.A tr2

\implies

\exists sv1 trv1 sv2 trv2.

istateV sv1 \wedge istateV sv2 \wedge corrState sv1 s1 \wedge corrState sv2 s2 \wedge

ψ s1 tr1 s2 tr2 Eq sv1 trv1 sv2 trv2

shows rsecure

unfolding rsecure-def2 **apply** safe

subgoal for s1 tr1 s2 tr2

using assms[of s1 tr1 s2 tr2]

using ψ -completedFrom ψ -def completedFromO-lastt **apply** clarsimp **by** metis .

proposition unwindCond-ex- ψ :

assumes unwind: unwindCond Δ

and Δ : Δ w w1 w2 s1 s2 statA sv1 sv2 statO **and** stat: (statA = Diff \longrightarrow statO = Diff)

and v: Opt.validFromS s1 tr1 Opt.completedFrom s1 tr1 Opt.validFromS s2 tr2 Opt.completedFrom s2 tr2

and tr14: Opt.A tr1 = Opt.A tr2

and r: reachO s1 reachO s2 reachV sv1 reachV sv2

shows \exists trv1 trv2. ψ s1 tr1 s2 tr2 statO sv1 trv1 sv2 trv2

using assms(2-)

proof(induction length tr1 + length tr2 w

arbitrary: w1 w2 s1 s2 statA sv1 sv2 statO tr1 tr2 rule: less2-induct')

case (less w tr1 tr2 w1 w2 s1 s2 statA sv1 sv2 statO)

note ok = \langle statA = Diff \longrightarrow statO = Diff \rangle

note Δ = \langle Δ w w1 w2 s1 s2 statA sv1 sv2 statO \rangle

note A34 = \langle Opt.A tr1 = Opt.A tr2 \rangle

note r34 = less.prem(8,9) **note** r12 = less.prem(10,11)

note r = r34 r12

note r3 = r34(1) **note** r4 = r34(2) **note** r1 = r12(1) **note** r2 = r12(2)

```

have i34: statA = Eq  $\longrightarrow$  isIntO s1 = isIntO s2
and f34: finalO s1 = finalO s2  $\wedge$  finalV sv1 = finalO s1  $\wedge$  finalV sv2 = finalO
s2
using  $\Delta$  unwind[unfolded unwindCond-def] r by auto

have proact-match:  $(\exists v < w. \text{proact } \Delta v w1 w2 s1 s2 \text{ statA sv1 sv2 statO}) \vee \text{react}$ 
 $\Delta w1 w2 s1 s2 \text{ statA sv1 sv2 statO}$ 
using  $\Delta$  unwind[unfolded unwindCond-def] r by auto
show ?case using proact-match proof safe
fix v assume v: v < w
assume proact  $\Delta v w1 w2 s1 s2 \text{ statA sv1 sv2 statO}$ 
thus ?thesis unfolding proact-def proof safe
assume sv1:  $\neg$  isSecV sv1  $\neg$  isIntV sv1 and move-1  $\Delta v w1 w2 s1 s2 \text{ statA}$ 
sv1 sv2 statO
then obtain sv1'
where 0: validTransV (sv1,sv1')
and  $\Delta$ :  $\Delta v w1 w2 s1 s2 \text{ statA sv1' sv2 statO}$ 
unfolding move-1-def by auto
have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain trv1 trv2 where  $\psi$ :  $\psi s1 tr1 s2 tr2 \text{ statO sv1' trv1 sv2 trv2}$ 
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2 statO, simplified,
OF  $\Delta ok$  - - - - r34 r1' r2]
using A34 less.premis(3-6) by blast
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using  $\psi ok 0 sv1$  unfolding  $\psi$ -def Van.completedFrom-def by auto
next
assume sv2:  $\neg$  isSecV sv2  $\neg$  isIntV sv2 and move-2  $\Delta v w1 w2 s1 s2 \text{ statA}$ 
sv1 sv2 statO
then obtain sv2'
where 0: validTransV (sv2,sv2')
and  $\Delta$ :  $\Delta v w1 w2 s1 s2 \text{ statA sv1 sv2' statO}$ 
unfolding move-2-def by auto
have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain trv1 trv2 where  $\psi$ :  $\psi s1 tr1 s2 tr2 \text{ statO sv1 trv1 sv2' trv2}$ 
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1 sv2' statO, simplified,
OF  $\Delta ok$  - - - - r34 r1 r2']
using A34 less.premis(3-6) by blast
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\psi ok 0 sv2$  unfolding  $\psi$ -def Van.completedFrom-def by auto
next
assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
and move-12  $\Delta v w1 w2 s1 s2 \text{ statA sv1 sv2 statO}$ 
then obtain sv1' sv2' statO'
where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1')  $\neg$  isSecV sv1
validTransV (sv2,sv2')  $\neg$  isSecV sv2
Van.eqAct sv1 sv2
and  $\Delta$ :  $\Delta v w1 w2 s1 s2 \text{ statA sv1' sv2' statO'}$ 
unfolding move-12-def by auto

```



```

    have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
    have ok': statA = Diff  $\longrightarrow$  statO' = Diff using ok 0 unfolding sstatO'-def
by (cases statO, auto)
    obtain trv1 trv2 where  $\psi: \psi\ s1\ tr1\ s2\ tr2\ statO'\ sv1'\ trv1\ sv2'\ trv2$ 
    using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2' statO', simplified,
OF  $\Delta\ ok'\ -\ -\ -\ -\ r34\ r12'$ ]
    using A34 less.prem(3-6) by blast
    show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
    using  $\psi\ ok'\ 0\ sv12$  unfolding  $\psi$ -def sstatO'-def Van.completedFrom-def
    using Van.A.Cons-unfold Van.eqAct-def completedFromO-lastt less.prem(4)
    less.prem(6) by auto
qed
next
assume m: react  $\Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
show ?thesis
proof(cases length tr1  $\leq$  Suc 0)
  case True note tr1 = True
  hence tr1 = []  $\vee$  tr1 = [s1]
  by (metis Simple-Transition-System.validFromS-Cons-iff Suc-length-conv le-Suc-eq
le-zero-eq length-0-conv less.prem(3))
  hence finalO s1 using less(3-6)
  using Opt.completed-Cons Opt.completed-Nil by blast
  hence f4: finalO s2 using f34 by blast
  hence tr2: tr2 = []  $\vee$  tr2 = [s2]
  by (metis Opt.final-def Simple-Transition-System.validFromS-Cons-iff less.prem(5)
neq-Nil-conv)
  show ?thesis apply(rule exI[of - [sv1]], rule exI[of - [sv2]]) using tr1 tr2
  using f4 f34
  using completedFromO-lastt less.prem(4)
  by (auto simp add: lastt-def  $\psi$ -def)
next
case False
then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1'  $\neq$  []
  by (cases tr1, auto)
  have s13[simp]: s13 = s1 using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
  by (simp add: Opt.validFromS-Cons-iff tr1)
  obtain s1' where
  trn3: validTransO (s1, s1') and
  tr1': Opt.validFromS s1' tr1' using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
  unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
  have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv
snd-conv)
  have f3:  $\neg$  finalO s1 using Opt.final-def trn3 by blast
  hence f4:  $\neg$  finalO s2 using f34 by blast
  hence tr2:  $\neg$  length tr2  $\leq$  Suc 0
  by (metis Opt.completed-Cons Simple-Transition-System.validFromS-Cons-iff

```

bot-nat-0.extremum completedFromO-def length-Cons less.premis(5) less.premis(6)
neq-Nil-conv not-less-eq-eq)

```

then obtain s24 tr2' where tr2: tr2 = s24 # tr2' and tr2'NE: tr2' ≠ []
by (cases tr2, auto)
have s24 [simp]: s24 = s2 using ⟨Opt.validFromS s2 tr2⟩
by (simp add: Opt.validFromS-Cons-iff tr2)
obtain s2' where
trn4: validTransO (s2,s2') ∨ (s2 = s2' ∧ tr2' = []) and
tr2': Opt.validFromS s2' tr2' using ⟨Opt.validFromS s2 tr2⟩
unfolding tr2 s24 using Opt.validFromS-Cons-iff by auto
have r34': reachO s1' reachO s2'
using r3 trn3 r4 trn4 by (metis Opt.reach.Step fst-conv snd-conv)+
note r3' = r34'(1) note r4' = r34'(2)
define statA' where statA': statA' = sstatA' statA s1 s2
have ¬ isIntO s1 ∨ ¬ isIntO s2 ∨ (isIntO s1 ∧ isIntO s2)
by auto
thus ?thesis
proof safe
assume isAO3: ¬ isIntO s1
have O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'
using isAO3 unfolding tr1 by auto
have A34': Opt.A tr1' = Opt.A tr2
using A34 O33'(2) by auto
have m: match1 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
have (∃ w1' < w1. ∃ w2' < w2. ¬ isSecO s1 ∧ Δ ∞ w1' w2' s1' s2 statA sv1
sv2 statO) ∨
(∃ w2' < w2. eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ match1-1 Δ ∞ w2' s1 s1'
s2 statA sv1 sv2 statO) ∨
(eqSec sv1 s1 ∧ ¬ isSecV sv2 ∧ Van.eqAct sv1 sv2 ∧ match1-12 Δ ∞
∞ s1 s1' s2 statA sv1 sv2 statO)
using m isAO3 trn3 ok unfolding match1-def by auto
thus ?thesis
proof safe
fix w1' w2'
assume ¬ isSecO s1 and Δ: Δ ∞ w1' w2' s1' s2 statA sv1 sv2 statO
hence S3: Opt.S tr1' = Opt.S tr1 unfolding tr1 by auto
obtain trv1 trv2 where ψ: ψ s1 tr1' s2 tr2 statO sv1 trv1 sv2 trv2
using less(1)[of tr1' tr2, OF - Δ - - - - - r3' r4 r12, unfolded O33',
simplified]
using less.premis tr1' ok A34' f3 f4 unfolding tr1 Opt.completedFrom-def
by (auto split: if-splits simp: ψ-def lastt-def)
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
using ψ O33' S3 unfolding ψ-def
using completedFromO-lastt less.premis(4)
by (auto simp add: tr1 tr1'NE)
next
fix w2'

```

```

assume trn13: eqSec sv1 s1 and
Atrn1:  $\neg$  isIntV sv1 and match1-1  $\Delta \infty w2' s1 s1' s2$  statA sv1 sv2 statO
then obtain sv1' where
  trn1: validTransV (sv1,sv1') and
   $\Delta$ :  $\Delta \infty \infty w2' s1' s2$  statA sv1' sv2 statO
unfolding match1-1-def by auto
  have r1': reachV sv1' using r1 trn1 by (metis Van.reach.Step fst-conv
snd-conv)
obtain trv1 trv2 where  $\psi$ :  $\psi s1 tr1' s2 tr2$  statO sv1' trv1 sv2 trv2
using less(1)[of tr1' tr2, OF -  $\Delta$  - - - - - r3' r4 r1' r2, unfolded O33',
simplified]
using less.premis tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def

by (auto simp:  $\psi$ -def lastt-def split: if-splits)
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
using  $\psi$  O33' unfolding tr1 tr2 Van.completedFrom-def
using Van.validFromS-Cons trn1 tr1'NE tr2'NE
using isAO3 ok Atrn1 eqSec-S-Cons trn13
unfolding  $\psi$ -def using completedFromO-lastt less.premis(4) tr1 by auto

next
assume sv2:  $\neg$  isSecV sv2 and trn13: eqSec sv1 s1 and
Atrn12: Van.eqAct sv1 sv2 and match1-12  $\Delta \infty \infty s1 s1' s2$  statA sv1
sv2 statO
then obtain sv1' sv2' statO' where
  statO': statO' = sstatO' statO sv1 sv2 and
  trn1: validTransV (sv1,sv1') and
  trn2: validTransV (sv2,sv2') and
   $\Delta$ :  $\Delta \infty \infty \infty s1' s2$  statA sv1' sv2' statO'
unfolding match1-12-def by auto
have r12': reachV sv1' reachV sv2'
using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
obtain trv1 trv2 where  $\psi$ :  $\psi s1' tr1' s2 tr2$  statO' sv1' trv1 sv2' trv2
using less(1)[of tr1' tr2, OF -  $\Delta$  - - - - - r3' r4 r12', unfolded O33',
simplified]
using less.premis tr1' ok A34' f3 f4 unfolding tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
by auto presburger+
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
using  $\psi$  O33' tr1'NE tr2'NE sv2
using Van.validFromS-Cons trn1 trn2
using isAO3 ok Atrn12 eqSec-S-Cons trn13 f3 f34 s13
unfolding  $\psi$ -def tr1 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def
using Van.A.Cons-unfold tr1' trn3 by auto
qed
next
assume isAO4:  $\neg$  isIntO s2
have O44': Opt.O tr2 = Opt.O tr2' Opt.A tr2 = Opt.A tr2'

```

```

using isAO4 unfolding tr2 by auto
have A34': Opt.A tr1 = Opt.A tr2'
using A34 O44'(2) by auto
have m: match2  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
have ( $\exists w1' < w1. \exists w2' < w2. \neg isSecO s2 \wedge \Delta \infty w1' w2' s1 s2' statA sv1$ 
sv2 statO)  $\vee$ 
( $\exists w1' < w1. eqSec sv2 s2 \wedge \neg isIntV sv2 \wedge match2-1 \Delta w1' \infty s1 s2$ 
s2' statA sv1 sv2 statO)  $\vee$ 
( $\neg isSecV sv1 \wedge eqSec sv2 s2 \wedge Van.eqAct sv1 sv2 \wedge match2-12 \Delta \infty$ 
 $\infty s1 s2 s2' statA sv1 sv2 statO$ )
using m isAO4 trn4 ok tr2'NE unfolding match2-def by auto
thus ?thesis
proof safe
fix w1' w2'
assume  $\neg isSecO s2$  and  $\Delta: \Delta \infty w1' w2' s1 s2' statA sv1 sv2 statO$ 
hence S4: Opt.S tr2' = Opt.S tr2 unfolding tr2 by auto
obtain trv1 trv2 where  $\psi: \psi s1 tr1 s2' tr2' statO sv1 trv1 sv2 trv2$ 
using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4', simplified]
using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
by (cases isIntO s2, auto)
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
using  $\psi$  O44' S4 unfolding  $\psi$ -def
using completedFromO-lastt less.premis(6)
unfolding Opt.completedFrom-def using tr2 tr2'NE by auto
next
fix w1'
assume trn24: eqSec sv2 s2 and
Atrn2:  $\neg isIntV sv2$  and match2-1  $\Delta w1' \infty s1 s2 s2' statA sv1 sv2 statO$ 
then obtain sv2' where trn2: validTransV (sv2,sv2') and
 $\Delta: \Delta \infty w1' \infty s1 s2' statA sv1 sv2' statO$ 
unfolding match2-1-def by auto
have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
obtain trv1 trv2 where  $\psi: \psi s1 tr1 s2' tr2' statO sv1 trv1 sv2' trv2$ 
using less(1)[of tr1 tr2', OF -  $\Delta$  - - - - - r3 r4' r1 r2', simplified]
using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
by (cases isIntO s2, auto)
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\psi$  tr1'NE tr2'NE
using Van.validFromS-Cons trn2
using isAO4 ok Atrn2 eqSec-S-Cons trn24
unfolding  $\psi$ -def tr1 tr2 s13 s24 Van.completedFrom-def lastt-def by auto
next
assume sv1:  $\neg isSecV sv1$  and trn24: eqSec sv2 s2 and
Atrn12: Van.eqAct sv1 sv2 and match2-12  $\Delta \infty \infty s1 s2 s2' statA sv1$ 
sv2 statO
then obtain sv1' sv2' statO' where
statO': statO' = sstatO' statO sv1 sv2 and

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```

    trn1: validTransV (sv1,sv1') and
    trn2: validTransV (sv2,sv2') and
    Δ: Δ ∞ ∞ ∞ s1 s2' statA sv1' sv2' statO'
    unfolding match2-12-def by auto
    have r12': reachV sv1' reachV sv2'
    using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
    obtain trv1 trv2 where ψ: ψ s1 tr1 s2' tr2' statO' sv1' trv1 sv2' trv2
    using less(1)[of tr1 tr2', OF - Δ - - - - - r3 r4' r12', simplified]
    using less.premis tr2' ok A34' tr1'NE tr2'NE unfolding tr1 tr2 Opt.completedFrom-def
statO' sstatO'-def
    by (cases isIntO s2, auto)
    show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
    using ψ O44' tr1'NE tr2'NE sv1
    using Van.validFromS-Cons trn1 trn2
    using isAO4 ok Atrn12 eqSec-S-Cons trn24
    unfolding ψ-def tr2 tr1'NE Van.completedFrom-def Van.eqAct-def
statO' sstatO'-def
    using Van.A.Cons-unfold tr2' trn4 by auto
qed
next
assume isAO34: isIntO s1 isIntO s2
have A34': getActO s1 = getActO s2 Opt.A tr1' = Opt.A tr2'
using A34 isAO34 tr1'NE tr2'NE unfolding tr1 tr2 by auto
have O33': Opt.O tr1 = getObsO s1 # Opt.O tr1' and
    O44': Opt.O tr2 = getObsO s2 # Opt.O tr2'
using isAO34 tr1'NE tr2'NE unfolding tr1 s13 tr2 s24 by auto
have m: match12 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
statA' react-def by auto
have trn34: getObsO s1 = getObsO s2 ∨ statA' = Diff
using isAO34 unfolding statA' sstatA'-def by (cases statA,auto)
have (∃ w1' < w1. ∃ w2' < w2. ¬ isSecO s1 ∧ ¬ isSecO s2 ∧ (statA = statA'
∨ statO = Diff) ∧
    Δ ∞ w1' w2' s1' s2' statA' sv1 sv2 statO)
    ∨
    (∃ w2' < w2. ¬ isSecO s2 ∧ eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧
    (statA = statA' ∨ statO = Diff) ∧
    match12-1 Δ ∞ w2' s1' s2' statA' sv1 sv2 statO)
    ∨
    (∃ w1' < w1. ¬ isSecO s1 ∧ eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧
    (statA = statA' ∨ statO = Diff) ∧
    match12-2 Δ w1' ∞ s1' s2' statA' sv1 sv2 statO)
    ∨
    (eqSec sv1 s1 ∧ eqSec sv2 s2 ∧ Van.eqAct sv1 sv2 ∧
    match12-12 Δ ∞ ∞ s1' s2' statA' sv1 sv2 statO)
(is ?K1 ∨ ?K2 ∨ ?K3 ∨ ?K4)
using m[unfolded match12-def, rule-format, of s1' s2']
isAO34 A34' trn3 trn4 tr1'NE tr2'NE
unfolding s13 s24 trn34 statA' Opt.eqAct-def sstatA'-def by auto

```

```

thus ?thesis proof (elim disjE)
  assume K1: ?K1
  then obtain w1' w2' where  $\Delta: \Delta \infty w1' w2' s1' s2' statA' sv1 sv2 statO$ 
by auto
  have ok': (statA' = Diff  $\longrightarrow$  statO = Diff)
  using ok K1 unfolding statA' using isAO34' by auto
  obtain trv1 trv2 where  $\psi: \psi s1' tr1' s2' tr2' statO sv1 trv1 sv2 trv2$ 
  using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r12, simplified]
  using less.premis tr1' tr2' ok' A34' isAO34' tr1'NE tr2'NE unfolding tr1
tr2 Opt.completedFrom-def by auto
  show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
  using  $\psi$  trn34 O33' O44' K1 ok unfolding  $\psi$ -def tr1 tr2
  using completedFromO-lastt less.premis(4,6)
  unfolding Opt.completedFrom-def using tr1 tr2 tr1'NE tr2'NE by auto
next
  assume K2: ?K2
  then obtain w2' sv1' where
  trn1: validTransV (sv1,sv1') and
  trn13: eqSec sv1 s1 and
  Atrn1:  $\neg$  isIntV sv1 and ok': (statA' = statA  $\vee$  statO = Diff) and
 $\Delta: \Delta \infty \infty w2' s1' s2' statA' sv1' sv2 statO$ 
  unfolding match12-1-def by auto
  have r1': reachV sv1' using r1 trn1 by (metis Van.reach.Step fst-conv
snd-conv)
  obtain trv1 trv2 where  $\psi: \psi s1' tr1' s2' tr2' statO sv1' trv1 sv2 trv2$ 
  using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r1' r2, simplified]
  using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
  show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])
  using  $\psi$  O33' O44' tr1'NE tr2'NE unfolding tr1 tr2
  using Van.validFromS-Cons trn1 ok
  using K2 ok' Atrn1 eqSec-S-Cons trn13 trn34
  unfolding statA' Van.completedFrom-def eqSec-def
  using s13 tr1 tr1' tr2' trn3 trn4
by simp (smt (verit, best) Opt.S.Cons-unfold Simple-Transition-System.lastt-Cons

Van.A.Cons-unfold Van.O.Cons-unfold  $\psi$ -def completedFromO-lastt f3 f34
lastt-Nil
less.premis(4) status.simps(1))
next
  assume K3: ?K3
  then obtain w1' sv2' where
  trn2: validTransV (sv2,sv2') and
  trn24: eqSec sv2 s2 and
  Atrn2:  $\neg$  isIntV sv2 and ok': (statA' = statA  $\vee$  statO = Diff) and
 $\Delta: \Delta \infty w1' \infty s1' s2' statA' sv1 sv2' statO$ 
  unfolding match12-2-def by auto
  have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)

```

```

obtain trv1 trv2 where  $\psi: \psi\ s1'\ tr1'\ s2'\ tr2'\ statO\ sv1\ trv1\ sv2'\ trv2$ 
using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r1 r2', simplified]
using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
show ?thesis apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\psi\ O33'\ O44'\ tr1'\ NE\ tr2'\ NE$  unfolding  $\psi$ -def tr1 tr2
using Van.validFromS-Cons trn2 ok
using K3 ok' Atrn2 eqSec-S-Cons trn24 trn34
unfolding statA' Van.completedFrom-def
using tr1' tr2' trn3 trn4 by force
next
assume K4: ?K4
then obtain sv1' sv2' statO' where 0:
  statO' = sstatO' statO sv1 sv2
  validTransV (sv1,sv1')
  eqSec sv1 s1
  validTransV (sv2,sv2')
  eqSec sv2 s2
  Van.eqAct sv1 sv2
  and ok': statA' = Diff  $\longrightarrow$  statO' = Diff and  $\Delta: \Delta \infty \infty \infty s1' s2'$ 
statA' sv1' sv2' statO'
unfolding match12-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0
by (metis Van.reach.Step fst-conv snd-conv)+
obtain trv1 trv2 where  $\psi: \psi\ s1'\ tr1'\ s2'\ tr2'\ statO'\ sv1'\ trv1\ sv2'\ trv2$ 
using less(1)[of tr1' tr2', OF -  $\Delta$  - - - - - r34' r12', simplified]
using less.premis tr1' tr2' ok' A34' tr1'NE tr2'NE unfolding tr1 tr2
Opt.completedFrom-def by auto
show ?thesis apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 #
trv2])
using trn34
using  $\psi\ O33'\ O44'\ isAO34\ tr1'\ NE\ tr2'\ NE$  unfolding  $\psi$ -def tr1 tr2
using Van.validFromS-Cons 0
using K4 eqSec-S-Cons
unfolding statA' Van.eqAct-def Van.completedFrom-def match12-12-def
sstatO'-def
by simp (smt (z3) Simple-Transition-System.lastt-Cons Van.A.Cons-unfold
Van.O.Cons-unfold list.inject status.exhaust status.simps(1) tr1' tr2' trn3 trn4
newStat.simps(4) newStat-diff)
qed
qed
qed
qed
qed

```

lemma *unwindCond-final:*

unwindCond $\Delta \implies reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $(finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO s2)$
unfolding *unwindCond-def*
unfolding *proact-def react-def match1-def match1-1-def*
by *auto*

definition $\varphi \Delta w w1 w2 w1' w2' statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2$
 $trv2 statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $(length\ trv1 > Suc\ 0 \vee w1' \leq w1) \wedge (length\ trv2 > Suc\ 0 \vee w2' \leq w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \longrightarrow (statOO = Diff \longleftrightarrow Van.O\ trv1 \neq Van.O\ trv2)) \wedge$
 $(statA = Eq \longrightarrow (statAA = Diff \longleftrightarrow Opt.O\ tr1 \neq Opt.O\ tr2)) \wedge$

 $(statO = Diff \longrightarrow statOO = Diff) \wedge$
 $(statAA = Diff \longrightarrow statOO = Diff) \wedge$
 $\Delta w w1' w2' (lastt\ s1\ tr1) (lastt\ s2\ tr2) statAA (lastt\ sv1\ trv1) (lastt\ sv2\ trv2)$
 $statOO$

lemma φ -*final*:

assumes *unw*: *unwindCond* Δ

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *vtr14*: *Opt.validFromS* *s1* *tr1* *Opt.validFromS* *s2* *tr2*

and φ : $\varphi \Delta w w1 w2 w1' w2' statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2 trv2$
 $statOO$

shows $(finalV (lastt\ sv1\ trv1) \longleftrightarrow finalO (lastt\ s1\ tr1)) \wedge (finalV (lastt\ sv2\ trv2)$
 $\longleftrightarrow finalO (lastt\ s2\ tr2))$

proof–

have *rsv12*: *Van.validFromS* *sv1* *trv1* \longrightarrow *reachV* $(lastt\ sv1\ trv1)$

Van.validFromS *sv2* *trv2* \longrightarrow *reachV* $(lastt\ sv2\ trv2)$ **using** *r*

by (*simp add*: *Van.reach-validFromS-reach lastt-def*)**+**

have *rs14*: *Opt.validFromS* *s1* *tr1* \longrightarrow *reachO* $(lastt\ s1\ tr1)$

Opt.validFromS *s2* *tr2* \longrightarrow *reachO* $(lastt\ s2\ tr2)$ **using** *r*

by (*simp add*: *Opt.reach-validFromS-reach lastt-def*)**+**

show *?thesis* **using** φ [*unfolded* φ -*def*] *rsv12* *rs14* **using** *unw*[*unfolded* *unwind-*
Cond-def, *rule-format*,

of lastt s1 tr1 lastt s2 tr2 lastt sv1 trv1 lastt sv2 trv2 w w1' w2' statAA statOO]

using *vtr14*(1) *vtr14*(2) **by** *auto*

qed

lemma φ -*completedFrom*: *unwindCond* $\Delta \implies$

reachO *s1* \implies *reachO* *s2* \implies *reachV* *sv1* \implies *reachV* *sv2* \implies

Opt.validFromS *s1* *tr1* \implies *completedFromO* *s1* *tr1* \implies

Opt.validFromS *s2* *tr2* \implies *completedFromO* *s2* *tr2* \implies

$\varphi \Delta \text{ statA } w \ w1 \ w2 \ w1' \ w2' \ s1 \ tr1 \ s2 \ tr2 \ \text{statAA} \ \text{statO} \ sv1 \ trv1 \ sv2 \ trv2 \ \text{statOO}$
 $\implies \text{completedFromV } sv1 \ trv1 \ \wedge \ \text{completedFromV } sv2 \ trv2$
using φ -final
by (*metis* *Van.completedFrom-def* *completedFromO-lastt* *lastt-def*)

lemma *unwindCond-ex- φ* :
assumes *unwind*: *unwindCond* Δ
and Δ : $\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *stat*: (*statA* = *Diff* \longrightarrow *statO* = *Diff*)
and *v*: *Opt.validFromS* *s1* *tr1* *Opt.validFromS* *s2* *tr2*
and *i*: *isIntO* (*lastt* *s1* *tr1*) *isIntO* (*lastt* *s2* *tr2*)
and *nev*: *never isIntO* (*butlast* *tr1*) *never isIntO* (*butlast* *tr2*)
shows $\exists w' \ w1' \ w2' \ trv1 \ trv2 \ \text{statAA} \ \text{statOO}. \ \varphi \Delta \ w' \ w1 \ w2 \ w1' \ w2' \ \text{statA} \ s1 \ tr1$
 $s2 \ tr2 \ \text{statAA} \ \text{statO} \ sv1 \ trv1 \ sv2 \ trv2 \ \text{statOO}$
using *assms*(2-)
proof(*induction* *length* *tr1* + *length* *tr2* *w*
arbitrary: *w1* *w2* *s1* *s2* *statA* *sv1* *sv2* *statO* *tr1* *tr2* *rule*: *less2-induct'*)
case (*less* *w* *tr1* *tr2* *w1* *w2* *s1* *s2* *statA* *sv1* *sv2* *statO*)
note *ok* = $\langle \text{statA} = \text{Diff} \longrightarrow \text{statO} = \text{Diff} \rangle$
note $\Delta = \langle \Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO} \rangle$
note *r34* = *less*(4,5) **note** *r12* = *less*(6,7)
note *r* = *r34* *r12*
note *r3* = *r34*(1) **note** *r4* = *r34*(2) **note** *r1* = *r12*(1) **note** *r2* = *r12*(2)
note *nev34* = *less*(13,14)
note *nev3* = *nev34*(1) **note** *nev4* = *nev34*(2)

have *i34*: *statA* = *Eq* \longrightarrow *isIntO* *s1* = *isIntO* *s2*
and *f34*: *finalO* *s1* = *finalO* *s2* \wedge *finalV* *sv1* = *finalO* *s1* \wedge *finalV* *sv2* = *finalO* *s2*
using Δ *unwind*[*unfolded* *unwindCond-def*] *r* **by** *auto*

note *is1* = $\langle \text{isIntO} \ (\text{lastt} \ s1 \ tr1) \rangle$
note *is2* = $\langle \text{isIntO} \ (\text{lastt} \ s2 \ tr2) \rangle$
note *utr1* = $\langle \text{Opt.validFromS} \ s1 \ tr1 \rangle$
note *utr2* = $\langle \text{Opt.validFromS} \ s2 \ tr2 \rangle$

have *proact-match*: ($\exists v < w. \ \text{proact} \ \Delta \ v \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$) \vee *react* $\Delta \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$
using Δ *unwind*[*unfolded* *unwindCond-def*] *r* **by** *auto*
show *?case* **using** *proact-match* **proof** *safe*
fix *v* **assume** *v*: *v* < *w*
assume *proact* $\Delta \ v \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1 \ sv2 \ \text{statO}$
thus *?thesis* **unfolding** *proact-def* **proof** *safe*
assume *sv1*: $\neg \text{isSecV} \ sv1 \ \neg \text{isIntV} \ sv1$ **and** *move-1* $\Delta \ v \ w1 \ w2 \ s1 \ s2 \ \text{statA}$
 $sv1 \ sv2 \ \text{statO}$
then obtain *sv1'*
where *0*: *validTransV* (*sv1*, *sv1'*)
and Δ : $\Delta \ v \ w1 \ w2 \ s1 \ s2 \ \text{statA} \ sv1' \ sv2 \ \text{statO}$

```

unfolding move-1-def by auto
have r1': reachV sv1' using r1 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w' w1' w2' trv1 trv2 statAA statOO where  $\varphi: \varphi \Delta w' w1 w2 w1' w2'$ 
statA s1 tr1 s2 tr2 statAA statO sv1' trv1 sv2 trv2 statOO
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2 statO, simplified,
OF  $\Delta r34 r1' r2 ok$ ]
using is1 is2 nev3 nev4 vtr1 vtr2 by blast
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1'])
apply(rule exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of
- trv2])
using  $\varphi ok 0 sv1$  unfolding  $\varphi$ -def by auto
next
assume sv2:  $\neg isSecV sv2 \neg isIntV sv2$  and move-2  $\Delta v w1 w2 s1 s2 statA$ 
sv1 sv2 statO
then obtain sv2'
where 0: validTransV (sv2,sv2')
and  $\Delta: \Delta v w1 w2 s1 s2 statA sv1 sv2' statO$ 
unfolding move-2-def by auto
have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w' w1' w2' trv1 trv2 statAA statOO where  $\varphi: \varphi \Delta w' w1 w2 w1' w2'$ 
statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2' trv2 statOO
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1 sv2' statO, simplified,
OF  $\Delta r34 r1 r2' ok$ ]
using is1 is2 nev3 nev4 vtr1 vtr2 by blast
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\varphi ok 0 sv2$  unfolding  $\varphi$ -def by auto
next
assume sv12:  $\neg isSecV sv1 \neg isSecV sv2 Van.eqAct sv1 sv2$ 
and move-12  $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$ 
then obtain sv1' sv2' statO'
where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1')  $\neg isSecV sv1$ 
validTransV (sv2,sv2')  $\neg isSecV sv2$ 
Van.eqAct sv1 sv2
and  $\Delta: \Delta v w1 w2 s1 s2 statA sv1' sv2' statO'$ 
unfolding move-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
have ok': statA = Diff  $\longrightarrow$  statO' = Diff
using ok 0 unfolding sstatO'-def by (cases statO, auto)
obtain w' w1' w2' trv1 trv2 statAA statOO where  $\varphi: \varphi \Delta w' w1 w2 w1' w2'$ 
statA s1 tr1 s2 tr2 statAA statO' sv1' trv1 sv2' trv2 statOO
using less(2)[OF v, of tr1 tr2 w1 w2 s1 s2 statA sv1' sv2' statO', simplified,
OF  $\Delta r34 r12' ok$ ]
using is1 is2 nev3 nev4 vtr1 vtr2 by blast
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2'])

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apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
apply(rule exI[of - statAA]) apply(rule exI[of - statOO])
using  $\varphi$  ok' 0 sv12 nev unfolding  $\varphi$ -def sstatO'-def
by simp (smt (verit, ccfv-SIG) Statewise-Attacker-Mod.eqAct-def
  Van.A.Cons-unfold Van.O.Cons-unfold Van.Statewise-Attacker-Mod-axioms
  Van.validFromS-Cons list.inject newStat.simps(1) newStat.simps(4))
qed
next
assume m: react  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
define statA' where statA': statA' = sstatA' statA s1 s2
show ?thesis
proof(cases length tr1  $\leq$  Suc 0)
  case True
    hence tr1e: tr1 = []  $\vee$  tr1 = [s1]
    by (metis Opt.validFromS-singl-iff Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
  vtr1)
    hence Opt.A tr1 = [] by (simp add: True)
    hence Opt.A tr2 = [] using Opt.A.eq-Nil-iff nev4 by blast
    show ?thesis
    proof(cases length tr2  $\leq$  Suc 0)
      case True
        hence tr2e: tr2 = []  $\vee$  tr2 = [s2]
        by (metis Opt.validFromS-def Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
  list.sel(1) vtr2)
        show ?thesis apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule
  exI[of - w2])
        apply(rule exI[of - [sv1]], rule exI[of - [sv2]], rule exI[of - statA], rule exI[of
  - statO])
        using tr1e tr2e
        using f34  $\Delta$  apply (clarsimp simp:  $\varphi$ -def lastt-def)
        apply(cases statA, simp-all)
        apply (metis Opt.O.simps(4) Opt.S.simps(4) last-ConsL)
        by (metis Opt.S.simps(4) last.simps ok)
      next
        case False
          then obtain s24 tr2' where tr2: tr2 = s24 # tr2' and tr2'NE: tr2'  $\neq$  []
          by (cases tr2, auto)
          have s24[simp]: s24 = s2 using  $\langle$ Opt.validFromS s2 tr2 $\rangle$ 
          by (simp add: Opt.validFromS-Cons-iff tr2)
          obtain s2' where
            trn4: validTransO (s2,s2')  $\vee$  (s2 = s2'  $\wedge$  tr2' = []) and
            tr2': Opt.validFromS s2' tr2' using  $\langle$ Opt.validFromS s2 tr2 $\rangle$ 
          unfolding tr2 s24 using Opt.validFromS-Cons-iff by auto
          have r4': reachO s2'
          using r4 trn4 by (metis Opt.reach.Step fst-conv snd-conv)+
          have nev4': never isIntO (butlast tr2')
          by (metis Opt.O.Nil-iff Opt.O.eq-Nil-iff nev4 tr2)
          have isAO4:  $\neg$  isIntO s2
          using  $\langle$ Opt.A tr2 = [] $\rangle$  tr2 tr2'NE by auto

```

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have  $O_{44}'$ :  $Opt.O\ tr2 = Opt.O\ tr2' \ Opt.A\ tr2 = Opt.A\ tr2'$ 
using  $isAO_4 \langle Opt.A\ tr2 = [] \rangle\ tr2$  by auto
have  $m$ :  $match2\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$  using  $m$  unfolding
react-def by auto
have  $(\exists w1' < w1. \exists w2' < w2. \neg isSecO\ s2 \wedge \Delta \infty w1'\ w2'\ s1\ s2'\ statA\ sv1\ sv2\ statO) \vee$ 
 $(\exists w1' < w1. eqSec\ sv2\ s2 \wedge \neg isIntV\ sv2 \wedge match2-1\ \Delta\ w1' \infty s1\ s2\ s2'\ statA\ sv1\ sv2\ statO) \vee$ 
 $(\neg isSecV\ sv1 \wedge eqSec\ sv2\ s2 \wedge Van.eqAct\ sv1\ sv2 \wedge match2-12\ \Delta \infty \infty s1\ s2\ s2'\ statA\ sv1\ sv2\ statO)$ 
using  $isAO_4\ trn_4\ ok\ tr2'NE$ 
using  $m[unfolding\ match2-def, rule-format, of\ s2']$  by auto
thus ?thesis
proof safe
fix  $w1''\ w2''$  assume  $w12'$ :  $w1'' < w1\ w2'' < w2$ 
assume  $\neg isSecO\ s2$  and  $\Delta$ :  $\Delta \infty w1''\ w2''\ s1\ s2'\ statA\ sv1\ sv2\ statO$ 
hence  $S_4$ :  $Opt.S\ tr2' = Opt.S\ tr2$  unfolding  $tr2$  by auto
obtain  $w'\ w1'\ w2'\ trv1\ trv2\ statAA\ statOO$  where  $\varphi$ :  $\varphi\ \Delta\ w'\ w1''\ w2''\ w1'\ w2'\ statA\ s1\ tr1\ s2'\ tr2'\ statAA\ statO\ sv1\ trv1\ sv2\ trv2\ statOO$ 
using  $less(1)[of\ tr1\ tr2', OF - \Delta\ r_3\ r_4' \dots nev3\ nev4', unfolded\ tr2, simplified]$ 
using  $is1\ is2\ vtr1\ vtr2\ tr2'\ ok\ tr2'NE\ trn_4\ r1\ r2\ tr2$  by auto
show ?thesis apply( $rule\ exI[of - w']$ ) apply( $rule\ exI[of - w1']$ ) apply( $rule\ exI[of - w2']$ ) apply( $rule\ exI[of - trv1]$ ) apply( $rule\ exI[of - trv2]$ )
using  $\varphi\ O_{44}'\ S_4\ tr2\ tr2'NE\ trn_4\ tr2'\ w12'$  unfolding  $\varphi-def$  by auto
next
fix  $w1''$  assume  $w1'$ :  $w1'' < w1$ 
assume  $trn2_4$ :  $eqSec\ sv2\ s2$  and
 $Atrn2$ :  $\neg isIntV\ sv2$  and  $match2-1\ \Delta\ w1'' \infty s1\ s2\ s2'\ statA\ sv1\ sv2\ statO$ 
then obtain  $sv2'$  where  $trn2$ :  $validTransV\ (sv2, sv2')$  and
 $\Delta$ :  $\Delta \infty w1'' \infty s1\ s2'\ statA\ sv1\ sv2'\ statO$ 
unfolding  $match2-1-def$  by auto
have  $r2'$ :  $reachV\ sv2'$  using  $r2\ trn2$  by ( $metis\ Van.reach.Step\ fst-conv\ snd-conv$ )
obtain  $w'\ w1'\ w2'\ trv1\ trv2\ statAA\ statOO$  where  $\varphi$ :  $\varphi\ \Delta\ w'\ w1'' \infty w1'\ w2'\ statA\ s1\ tr1\ s2'\ tr2'\ statAA\ statO\ sv1\ trv1\ sv2'\ trv2\ statOO$ 
using  $less(1)[of\ tr1\ tr2', OF - \Delta\ r_3\ r_4'\ r1\ r2' \dots nev3\ nev4', unfolded\ tr2, simplified]$ 
using  $is1\ is2\ tr2'\ tr2\ vtr1\ ok\ tr2'NE\ trn_4$  by auto
show ?thesis apply( $rule\ exI[of - w']$ ) apply( $rule\ exI[of - w1']$ ) apply( $rule\ exI[of - w2']$ ) apply( $rule\ exI[of - trv1]$ ) apply( $rule\ exI[of - sv2 \# trv2]$ )
using  $\varphi\ tr2'NE$ 
using  $Van.validFromS-Cons\ trn2$ 
using  $isAO_4\ ok\ Atrn2\ eqSec-S-Cons\ trn2_4\ tr2'\ trn_4\ w1'$ 
unfolding  $\varphi-def\ tr2\ s2_4$ 
by auto
next
assume  $sv1$ :  $\neg isSecV\ sv1$  and  $trn2_4$ :  $eqSec\ sv2\ s2$  and
 $Atrn12$ :  $Van.eqAct\ sv1\ sv2$  and  $match2-12\ \Delta \infty \infty s1\ s2\ s2'\ statA\ sv1\ sv2$ 

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```

statO
  then obtain sv1' sv2' statO' where
    statO': statO' = sstatO' statO sv1 sv2 and
    trn1: validTransV (sv1,sv1') and
    trn2: validTransV (sv2,sv2') and
    Δ: Δ ∞ ∞ ∞ s1 s2' statA sv1' sv2' statO'
  unfolding match2-12-def by auto
  have r12': reachV sv1' reachV sv2'
  using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
  obtain w' w1' w2' trv1 trv2 statAA statOO where φ: φ Δ w' ∞ ∞ w1'
w2' statA s1 tr1 s2' tr2' statAA statO' sv1' trv1 sv2' trv2 statOO
  using less(1)[of tr1 tr2', OF - Δ r3 r4' r12' - - - - nev3 nev4', simplified]
  using is1 is2 vtr1 tr2 tr2' ok tr2'NE trn4 unfolding tr2 statO' sstatO'-def
by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
  using φ O44' tr2'NE sv1
  using Van.validFromS-Cons trn1 trn2
  using isAO4 ok Atrn12 eqSec-S-Cons trn24 tr2' trn4
  unfolding φ-def tr2 Van.completedFrom-def Van.eqAct-def statO' sstatO'-def

  by simp (smt (verit, ccfv-threshold) Van.A.Cons-unfold i34 is1 last-ConsL
lastt-def status.exhaust tr1e newStat.simps(2))
qed
qed
next
case False
then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1' ≠ []
  by (cases tr1, auto)
have s13[simp]: s13 = s1 using ⟨Opt.validFromS s1 tr1⟩
  by (simp add: Opt.validFromS-Cons-iff tr1)
obtain s1' where
  trn3: validTransO (s1,s1') and
  tr1': Opt.validFromS s1' tr1' using ⟨Opt.validFromS s1 tr1⟩
unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv snd-conv)
have f3: ¬ finalO s1 using Opt.final-def trn3 by blast
hence f4: ¬ finalO s2 using f34 by blast
have nev3': never isIntO (butlast tr1')
using nev3 tr1 tr1'NE by auto
have isAO3: ¬ isIntO s1 using less.prem(11) tr1 tr1'NE by auto
have O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'
using isAO3 unfolding tr1 by auto
  have m: match1 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
  have (∃ w1' < w1. ∃ w2' < w2. ¬ isSecO s1 ∧ Δ ∞ w1' w2' s1' s2 statA sv1
sv2 statO) ∨
    (∃ w2' < w2. eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ match1-1 Δ ∞ w2' s1 s1' s2
statA sv1 sv2 statO) ∨

```

$(eqSec\ sv1\ s1 \wedge \neg isSecV\ sv2 \wedge Van.eqAct\ sv1\ sv2 \wedge match1-12\ \Delta\ \infty\ \infty$
 $s1\ s1'\ s2\ statA\ sv1\ sv2\ statO)$
using $m\ isAO3\ trn3\ ok$ **unfolding** $match1-def$ **by** $auto$
thus $?thesis$
proof $safe$
fix $w1''\ w2''$ **assume** $w12'$: $w1'' < w1\ w2'' < w2$
assume $\neg isSecO\ s1$ **and** Δ : $\Delta\ \infty\ w1''\ w2''\ s1'\ s2\ statA\ sv1\ sv2\ statO$
hence $S3$: $Opt.S\ tr1' = Opt.S\ tr1$ **unfolding** $tr1$ **by** $auto$
obtain $w'\ w1'\ w2'\ trv1\ trv2\ statAA\ statOO$ **where** φ : $\varphi\ \Delta\ w'\ w1''\ w2''\ w1'$
 $w2'\ statA\ s1'\ tr1'\ s2\ tr2\ statAA\ statO\ sv1\ trv1\ sv2\ trv2\ statOO$
using $less(1)[of\ tr1'\ tr2,\ OF - \Delta\ r3'\ r4\ r12,\ unfolded\ O33',\ simplified]$
using $is1\ is2\ tr1'\ ok\ f3\ f4\ tr1'NE\ trn3\ O33'(1)\ nev3'\ nev4\ vtr2$ **unfolding**
 $tr1$ **by** $auto$
show $?thesis$ **apply**($rule\ exI[of - w']$) **apply**($rule\ exI[of - w1']$) **apply**($rule$
 $exI[of - w2']$) **apply**($rule\ exI[of - trv1]$) **apply**($rule\ exI[of - trv2]$)
using $\varphi\ O33'\ S3\ tr1\ tr1'NE\ tr1'\ trn3\ w12'$ **unfolding** $\varphi-def$ **by** $auto$
next
fix $w2''$ **assume** $w2'$: $w2'' < w2$
assume $trn13$: $eqSec\ sv1\ s1$ **and**
 $Atrn1$: $\neg isIntV\ sv1$ **and** $match1-1\ \Delta\ \infty\ w2''\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO$
then obtain $sv1'$ **where**
 $trn1$: $validTransV\ (sv1,sv1')$ **and**
 Δ : $\Delta\ \infty\ \infty\ w2''\ s1'\ s2\ statA\ sv1'\ sv2\ statO$
unfolding $match1-1-def$ **by** $auto$
have $r1'$: $reachV\ sv1'$ **using** $r1\ trn1$ **by** ($metis\ Van.reach.Step\ fst-conv$
 $snd-conv$)
obtain $w'\ w1'\ w2'\ trv1\ trv2\ statAA\ statOO$ **where** φ : $\varphi\ \Delta\ w'\ \infty\ w2''\ w1'$
 $w2'\ statA\ s1'\ tr1'\ s2\ tr2\ statAA\ statO\ sv1'\ trv1\ sv2\ trv2\ statOO$
using $less(1)[of\ tr1'\ tr2,\ OF - \Delta\ r3'\ r4\ r1'\ r2,\ unfolded\ O33',\ simplified]$
using $is1\ is2\ tr1\ nev3'\ nev4\ vtr1\ vtr2\ tr1'\ ok\ f3\ f4\ tr1'NE\ trn3\ O33'(1)$
unfolding $tr1$ **by** $auto$
show $?thesis$ **apply**($rule\ exI[of - w']$) **apply**($rule\ exI[of - w1']$) **apply**($rule$
 $exI[of - w2']$) **apply**($rule\ exI[of - sv1\ \# trv1]$) **apply**($rule\ exI[of - trv2]$)
using $\varphi\ O33'$ **unfolding** $\varphi-def\ tr1\ Van.completedFrom-def$
using $Van.validFromS-Cons\ trn1\ tr1'NE\ tr1'\ trn3$
using $isAO3\ ok\ Atrn1\ eqSec-S-Cons\ trn13\ w2'$
by $auto$
next
assume $sv2$: $\neg isSecV\ sv2$ **and** $trn13$: $eqSec\ sv1\ s1$ **and**
 $Atrn12$: $Van.eqAct\ sv1\ sv2$ **and** $match1-12\ \Delta\ \infty\ \infty\ s1\ s1'\ s2\ statA\ sv1\ sv2$
 $statO$
then obtain $sv1'\ sv2'\ statO'$ **where**
 $statO'$: $statO' = sstatO'\ statO\ sv1\ sv2$ **and**
 $trn1$: $validTransV\ (sv1,sv1')$ **and**
 $trn2$: $validTransV\ (sv2,sv2')$ **and**
 Δ : $\Delta\ \infty\ \infty\ \infty\ s1'\ s2\ statA\ sv1'\ sv2'\ statO'$
unfolding $match1-12-def$ **by** $auto$
have $r12'$: $reachV\ sv1'\ reachV\ sv2'$
using $r1\ trn1\ r2\ trn2$ **by** ($metis\ Van.reach.Step\ fst-conv\ snd-conv$)+

obtain $w' w1' w2' trv1 trv2 statAA statOO$ **where** $\varphi: \varphi \Delta w' \infty \infty w1'$
 $w2' statA s1' tr1' s2 tr2 statAA statO' sv1' trv1 sv2' trv2 statOO$
using $less(1)[of\ tr1'\ tr2, OF - \Delta\ r3'\ r4\ r12', unfolded\ O33', simplified]$
using $less.premis\ tr1'\ ok\ f3\ f4\ tr1'NE\ trn3\ O33'(1)$ **unfolding** $tr1\ statO'$
 $sstatO'-def$ **by** *auto*

have $trv1NE: trv1 \neq []$ **and** $trv2NE: trv2 \neq []$ **using** φ **unfolding** $\varphi-def$ **by**
auto

have $[simp]: Van.O\ (sv1 \# trv1) = Van.O\ (sv2 \# trv2) \longleftrightarrow (isIntV\ sv1$
 $\longrightarrow getObsV\ sv1 = getObsV\ sv2) \wedge Van.O\ trv1 = Van.O\ trv2$
using $Atrn12\ trv1NE\ trv2NE$ **unfolding** $Van.O.map-filter\ Van.eqAct-def$ **by**
 $simp$

show $?thesis$ **apply**($rule\ exI[of - w']$) **apply**($rule\ exI[of - w1']$) **apply**($rule$
 $exI[of - w2']$) **apply**($rule\ exI[of - sv1 \# trv1]$) **apply**($rule\ exI[of - sv2 \# trv2]$)
using $\varphi\ O33'\ tr1'NE\ sv2$
using $Van.validFromS-Cons\ trn1\ trn2$
using $isAO3\ ok\ Atrn12\ eqSec-S-Cons\ trn13\ f3\ f34\ s13\ tr1'\ trn3$
unfolding $\varphi-def\ tr1\ Van.completedFrom-def\ Van.eqAct-def\ statO'\ sstatO'-def$
apply $clarsimp$

by ($smt\ (verit, ccfv-SIG)\ Van.A.Cons-unfold\ newStat.simps(1)\ newStat.simps(2)$
 $newStat.simps(4)$)
qed
qed
qed
qed

definition $\varphi a \Delta w w1 w2 w1' w2' statA s1 tr1 s2 tr2 statAA statO sv1 trv1 sv2$
 $trv2 statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge$
 $(length\ trv1 > Suc\ 0 \vee w1' < w1) \wedge (length\ trv2 > Suc\ 0 \vee w2' < w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge Van.S\ trv2 = Opt.S\ tr2 \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $(statO = Eq \longrightarrow (statOO = Diff \longleftrightarrow Van.O\ trv1 \neq Van.O\ trv2)) \wedge$
 $(statA = Eq \longrightarrow (statAA = Diff \longleftrightarrow Opt.O\ tr1 \neq Opt.O\ tr2)) \wedge$

 $(statO = Diff \longrightarrow statOO = Diff) \wedge$
 $(statAA = Diff \longrightarrow statOO = Diff) \wedge$
 $\Delta w w1' w2' (lastt\ s1\ tr1) (lastt\ s2\ tr2) statAA (lastt\ sv1\ trv1) (lastt\ sv2\ trv2)$
 $statOO$

lemma $unwindCond-ex-\varphi a-getActO:$
assumes $unwind: unwindCond\ \Delta$
and $\Delta: \Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and $r34: reachO\ s1\ reachO\ s2$ **and** $r12: reachV\ sv1\ reachV\ sv2$
and $stat: (statA = Diff \longrightarrow statO = Diff)$
and $v: validTransO\ (s1, s1')\ validTransO\ (s2, s2')$
and $i34: isIntO\ s1\ isIntO\ s2\ getActO\ s1 = getActO\ s2$

shows $\exists w1' w2' trv1 trv2 statOO$.
 $\varphi a \Delta \infty w1 w2 w1' w2' statA s1 [s1, s1'] s2 [s2, s2'] (sstatA' statA s1 s2)$
 $statO sv1 trv1 sv2 trv2 statOO$
using $\Delta r12 stat$
proof(*induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct*)
case (*less w w1 w2 sv1 sv2 statO*)
note $\Delta = \langle \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \rangle$
note $r12 = less.prem(2,3)$
note $r1 = r12(1)$ **note** $r2 = r12(2)$
note $r = r34 r12$
note $stat = \langle statA = Diff \longrightarrow statO = Diff \rangle$

have $f34: finalO s1 = finalO s2 \wedge finalV sv1 = finalO s1 \wedge finalV sv2 = finalO s2$
using $\Delta unwind[unfolded unwindCond-def] r$ **by** *auto*

have *proact-match*: $(\exists v < w. proact \Delta v w1 w2 s1 s2 statA sv1 sv2 statO) \vee react \Delta w1 w2 s1 s2 statA sv1 sv2 statO$
using $\Delta unwind[unfolded unwindCond-def] r$ **by** *auto*
show *?case using proact-match proof safe*
fix v **assume** $v: v < w$
assume *proact* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
thus *?thesis unfolding proact-def proof safe*
assume $sv1: \neg isSecV sv1 \neg isIntV sv1$ **and** *move-1* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
then obtain $sv1'$
where $0: validTransV (sv1, sv1')$
and $\Delta: \Delta v w1 w2 s1 s2 statA sv1' sv2 statO$
unfolding *move-1-def* **by** *auto*
have $r1': reachV sv1'$ **using** $r1 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w1' w2' trv1 trv2 statOO$ **where**
 $\varphi: \varphi a \Delta \infty w1 w2 w1' w2' statA s1 [s1, s1'] s2 [s2, s2'] (sstatA' statA s1 s2)$
 $statO sv1' trv1 sv2 trv2 statOO$
using *less(1)[OF v Δ r1' r2 stat]* **by** *auto*
show *?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])*
using $\varphi 0 sv1$ **unfolding** *φa-def* **apply** *simp*
by (*metis Van.validFromS-Cons*)
next
assume $sv2: \neg isSecV sv2 \neg isIntV sv2$ **and** *move-2* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
then obtain $sv2'$
where $0: validTransV (sv2, sv2')$
and $\Delta: \Delta v w1 w2 s1 s2 statA sv1 sv2' statO$
unfolding *move-2-def* **by** *auto*
have $r2': reachV sv2'$ **using** $r2 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w1' w2' trv1 trv2 statOO$ **where**
 $\varphi: \varphi a \Delta \infty w1 w2 w1' w2' statA s1 [s1, s1'] s2 [s2, s2'] (sstatA' statA s1 s2)$
 $statO sv1 trv1 sv2' trv2 statOO$


```

using less(1)[OF v Δ r1 r2' stat] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using φ 0 sv2 unfolding φa-def apply simp by (metis Van.validFromS-Cons)
next
assume sv12: ¬ isSecV sv1 ¬ isSecV sv2 Van.eqAct sv1 sv2
and move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO
then obtain sv1' sv2' statO'
where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1') ¬ isSecV sv1
validTransV (sv2,sv2') ¬ isSecV sv2
Van.eqAct sv1 sv2
and Δ: Δ v w1 w2 s1 s2 statA sv1' sv2' statO'
unfolding move-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+
have stat': statA = Diff ⟶ statO' = Diff
using stat 0 unfolding sstatO'-def by (cases statO, auto)
obtain w1' w2' trv1 trv2 statOO where
φ: φa Δ ∞ w1 w2 w1' w2' statA s1 [s1, s1'] s2 [s2, s2'] (sstatA' statA s1 s2)
statO' sv1' trv1 sv2' trv2 statOO
using less(1)[OF v Δ r12' stat'] unfolding φa-def apply simp by metis
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
using φ 0 unfolding φa-def sstatO'-def apply clarsimp apply(intro conjI)
subgoal by auto
subgoal by auto
subgoal by (metis Van.A.Cons-unfold Van.eqAct-def)
subgoal apply(rule exI[of - statOO]) apply simp
by (smt (verit, ccfv-threshold) Van.O.Cons-unfold Van.eqAct-def
list.inject newStat.simps(1) newStat.simps(3)) .
qed
next
assume m: react Δ w1 w2 s1 s2 statA sv1 sv2 statO
define statA' where statA': statA' = sstatA' statA s1 s2
have m: match12 Δ w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
have (∃ w1' w2'. w1' < w1 ∧ w2' < w2 ∧ ¬ isSecO s1 ∧ ¬ isSecO s2 ∧ (statA
= statA' ∨ statO = Diff) ∧
Δ ∞ w1' w2' s1' s2' statA' sv1 sv2 statO)
∨
(∃ w2' < w2. ¬ isSecO s2 ∧
eqSec sv1 s1 ∧ ¬ isIntV sv1 ∧ (statA = statA' ∨ statO = Diff) ∧
match12-1 Δ ∞ w2' s1' s2' statA' sv1 sv2 statO)
∨
(∃ w1' < w1. ¬ isSecO s1 ∧
eqSec sv2 s2 ∧ ¬ isIntV sv2 ∧ (statA = statA' ∨ statO = Diff) ∧
match12-2 Δ w1' ∞ s1' s2' statA' sv1 sv2 statO)
∨

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      (eqSec sv1 s1 ∧ eqSec sv2 s2 ∧ Van.eqAct sv1 sv2 ∧
       match12-12 Δ ∞ ∞ s1' s2' statA' sv1 sv2 statO)
using m unfolding match12-def
by (simp add: Opt.eqAct-def i34(1) i34(2) i34(3) statA' v(1) v(2))
thus ?thesis
apply(elim disjE exE)
  subgoal for w1' w2' apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - [sv1]]) apply(rule exI[of - [sv2]])
  apply(rule exI[of - statO])
  using stat unfolding φa-def statA'
  by (auto simp add: i34(1) i34(2) sstatA'-def lastt-def)
  subgoal for w2' apply(rule exI[of - ∞]) apply(rule exI[of - w2'])
  unfolding match12-1-def apply(elim conjE exE) subgoal for sv1'
  apply(rule exI[of - [sv1,sv1']]) apply(rule exI[of - [sv2]])
  apply(rule exI[of - statO])
  using stat unfolding φa-def statA'
  by (auto simp add: i34(1) i34(2) sstatA'-def lastt-def) .
  subgoal for w1' apply(rule exI[of - w1']) apply(rule exI[of - ∞])
  unfolding match12-2-def apply(elim conjE exE) subgoal for sv2'
  apply(rule exI[of - [sv1]]) apply(rule exI[of - [sv2,sv2']])
  apply(rule exI[of - statO])
  using stat unfolding φa-def statA'
  by (auto simp add: i34(1) i34(2) sstatA'-def lastt-def) .
  subgoal unfolding match12-12-def apply(elim conjE exE) subgoal for sv1'
sv2'
  apply(rule exI[of - ∞]) apply(rule exI[of - ∞])
  apply(rule exI[of - [sv1,sv1']]) apply(rule exI[of - [sv2,sv2']])
  apply(rule exI[of - sstatO' statO sv1 sv2])
  using stat unfolding φa-def statA'
  by (auto simp add: i34 i34 sstatA'-def sstatO'-def lastt-def Van.eqAct-def) . .
qed
qed

lemma unwindCond-ex-φa'-aux:
assumes unwind: unwindCond Δ
and Δ: Δ w w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: (statA = Diff → statO = Diff)
and tr14NE: tr1 ≠ [] tr2 ≠ []
and v3': Opt.validFromS s1 (tr1 ## s1') and v4': Opt.validFromS s2 (tr2 ##
s2')
and i: isIntO (lastt s1 tr1) isIntO (lastt s2 tr2)
and A34: getActO (lastt s1 tr1) = getActO (lastt s2 tr2)
and nev: never isIntO (butlast tr1) never isIntO (butlast tr2)
shows ∃ w1' w2' trv1' trv2' statAA' statOO'.
  φa Δ ∞ w1 w2 w1' w2' statA s1 (tr1 ## s1') s2 (tr2 ## s2') statAA' statO
sv1 trv1' sv2 trv2' statOO'
proof -
  have v3: Opt.validFromS s1 tr1 and s13': validTransO (lastt s1 tr1, s1')

```

apply (*metis* $v3'$ *Opt.validFromS-def* *Opt.validS-append1* *Nil-is-append-conv* *hd-append2*)
by (*metis* *Opt.validFromS-def* *Opt.validS-validTrans* *append-is-Nil-conv* *lastt-def*
list.distinct(1) *list.sel(1)* *tr14NE(1)* $v3'$)
have $v4$: *Opt.validFromS* $s2$ $tr2$ **and** $s24'$: *validTransO* (*lastt* $s2$ $tr2, s2'$)
apply (*metis* $v4'$ *Opt.validFromS-def* *Opt.validS-append1* *Nil-is-append-conv* *hd-append2*)
by (*metis* *Opt.validFromS-def* *Opt.validS-validTrans* *append-is-Nil-conv* *lastt-def*
list.sel(1) *list.simps(3)* *tr14NE(2)* $v4'$)

obtain ww $ww1$ $ww2$ $trv1$ $trv2$ *statAA* *statOO* **where** φ : $\varphi \Delta$ ww $w1$ $w2$ $ww1$
 $ww2$ *statA* $s1$ $tr1$ $s2$ $tr2$ *statAA* *statO* $sv1$ $trv1$ $sv2$ $trv2$ *statOO*
using *unwindCond-ex- φ [OF unwind Δ r *stat* $v3$ $v4$ i *nev*]* **by** *auto*

have $trv12NE$: $trv1 \neq []$ $trv2 \neq []$ **using** φ **unfolding** φ -*def* **by** *auto*

define $ss1$ $ss2$ $ssv1$ $ssv2$ **where** $ss1$: $ss1 \equiv \text{lastt } s1 \ tr1$ **and** $ss2$: $ss2 \equiv \text{lastt } s2$
 $tr2$

and $ssv1$: $ssv1 \equiv \text{lastt } sv1 \ trv1$ **and** $ssv2$: $ssv2 \equiv \text{lastt } sv2 \ trv2$

have $ss1l$: $ss1 = \text{last } tr1$ **by** (*simp* *add*: *lastt-def* $ss1$ *tr14NE(1)*)
have $tr1l$: $tr1 = \text{butlast } tr1 @ [ss1]$ **by** (*simp* *add*: $ss1l$ *tr14NE(1)*)
have $ss2l$: $ss2 = \text{last } tr2$ **by** (*simp* *add*: *lastt-def* $ss2$ *tr14NE(2)*)
have $tr2l$: $tr2 = \text{butlast } tr2 @ [ss2]$ **by** (*simp* *add*: $ss2l$ *tr14NE(2)*)
have $ssv1l$: $ssv1 = \text{last } trv1$ **using** φ **unfolding** φ -*def* **by** (*metis* *lastt-def* $ssv1$)
have $trv1l$: $trv1 = \text{butlast } trv1 @ [ssv1]$ **by** (*simp* *add*: $ssv1l$ *trv12NE(1)*)
have $ssv2l$: $ssv2 = \text{last } trv2$ **using** φ **unfolding** φ -*def* **by** (*metis* *lastt-def* $ssv2$)
have $trv2l$: $trv2 = \text{butlast } trv2 @ [ssv2]$ **by** (*simp* *add*: $ssv2l$ *trv12NE(2)*)

have $iss14$ [*simp*]: *isIntO* $ss1$ *isIntO* $ss2$ **using** i **unfolding** $ss1$ $ss2$ **by** *auto*

have $giss14$ [*simp*]: *getActO* $ss1 = \text{getActO } ss2$

using *A34* $ss1$ $ss2$ **by** *fastforce*

have [*simp*]: *Opt.O* ($tr1 \ ## \ s1'$) = *Opt.O* $tr1 \ ## \ \text{getObsO } ss1$
by (*metis* *Opt.O-def* $\langle \text{isIntO } ss1 \rangle$ *holds-filtermap-RCons* *snoc-eq-iff-butlast* $tr1l$)
have [*simp*]: *Opt.O* ($tr2 \ ## \ s2'$) = *Opt.O* $tr2 \ ## \ \text{getObsO } ss2$
by (*metis* *Opt.O-def* $\langle \text{isIntO } ss2 \rangle$ *holds-filtermap-RCons* *snoc-eq-iff-butlast* $tr2l$)

have [*simp*]: *Opt.A* ($tr1 \ ## \ s1'$) = *Opt.A* $tr1 \ ## \ \text{getActO } ss1$
by (*metis* *Opt.A-def* $\langle \text{isIntO } ss1 \rangle$ *holds-filtermap-RCons* *snoc-eq-iff-butlast* $tr1l$)
have [*simp*]: *Opt.A* ($tr2 \ ## \ s2'$) = *Opt.A* $tr2 \ ## \ \text{getActO } ss2$
by (*metis* *Opt.A-def* $\langle \text{isIntO } ss2 \rangle$ *holds-filtermap-RCons* *snoc-eq-iff-butlast* $tr2l$)
have [*simp*]: *Opt.A* ($tr1 \ ## \ s1'$) = *Opt.A* ($tr2 \ ## \ s2'$) \longleftrightarrow *Opt.A* $tr1 = \text{Opt.A}$
 $tr2$ **by** *simp*

have rss : *reachO* $ss1$ *reachO* $ss2$ *reachV* $ssv1$ *reachV* $ssv2$
using *Opt.reach-validFromS-reach* r $ss1l$ *tr14NE(1)* $v3$ **apply** *blast*
using *Opt.reach-validFromS-reach* $r(2)$ $ss2l$ *tr14NE(2)* $v4$ **apply** *blast*
using *Van.reach-validFromS-reach* φ -*def* φ $r(3)$ $ssv1l$
apply (*smt* (*verit*, *del-insts*))
using *Van.reach-validFromS-reach* φ -*def* φ $r(4)$ $ssv2l$

apply (*smt* (*verit*, *del-insts*)) .

have *stat*: *statAA* = *Diff* \longrightarrow *statOO* = *Diff*
and Δ : Δ *ww* *ww1* *ww2* *ss1* *ss2* *statAA* *ssv1* *ssv2* *statOO*
using φ **unfolding** φ -def *ss1*[*symmetric*] *ss2*[*symmetric*] *ssv1*[*symmetric*] *ssv2*[*symmetric*]
by *auto*

note *vs13* = *s13*'[*unfolded ss1*[*symmetric*]] **note** *vs24* = *s24*'[*unfolded ss2*[*symmetric*]]
have \exists *w1'* *w2'* *trv1'* *trv2'* *statA'* *statO'*.
 φ Δ ∞ *ww1* *ww2* *w1'* *w2'* *statAA* *ss1* [*ss1*,*s1*'] *ss2* [*ss2*,*s2*'] (*sstatA'* *statAA* *ss1* *ss2*) *statOO* *ssv1* *trv1'* *ssv2* *trv2'* *statO'*
using *unwindCond-ex- φ -getActO*[*OF unwind* Δ *rss* *stat* *vs13* *vs24* *iss14* *giss14*]
by *blast*

then obtain *w1'* *w2'* *trv1'* *trv2'* *statA'* *statO'* **where**
 φ 1: φ Δ ∞ *ww1* *ww2* *w1'* *w2'* *statAA* *ss1* [*ss1*,*s1*'] *ss2* [*ss2*,*s2*'] *statA'* *statOO* *ssv1* *trv1'* *ssv2* *trv2'* *statO'* **by** *auto*

have *trv12'NE*: *trv1'* \neq [] \wedge *trv2'* \neq [] **using** φ 1 **unfolding** φ a-def **by** *auto*

have [*simp*]: *Van.O* (*butlast* *trv1* @ *trv1'*) = *Van.O* *trv1* @ *Van.O* *trv1'*
using *trv12'NE* **unfolding** φ -def *Van.O.map-filter* *Opt.O.map-filter* **apply**(*subst* *butlast-append*) **by** *simp*

have [*simp*]: *Van.O* (*butlast* *trv2* @ *trv2'*) = *Van.O* *trv2* @ *Van.O* *trv2'*
using *trv12'NE* **unfolding** φ -def *Van.O.map-filter* *Opt.O.map-filter* **apply**(*subst* *butlast-append*) **by** *simp*

have *Van.A* *trv1'* = *Van.A* *trv2'* **using** φ 1 **unfolding** φ a-def **by** *auto*
moreover **have** *length* (*Van.O* *trv1'*) = *length* (*Van.A* *trv1'*) \wedge *length* (*Van.O* *trv2'*) = *length* (*Van.A* *trv2'*)
unfolding *Van.A.map-filter* *Van.O.map-filter* **by** *auto*
ultimately **have** *length* (*Van.O* *trv1'*) = *length* (*Van.O* *trv2'*) **by** *auto*
hence [*simp*]: *Van.O* *trv1* @ *Van.O* *trv1'* = *Van.O* *trv2* @ *Van.O* *trv2'* \longleftrightarrow
Van.O *trv1* = *Van.O* *trv2* \wedge *Van.O* *trv1'* = *Van.O* *trv2'* **by** *auto*

have *len*: *trv1* \neq [] \wedge *trv2* \neq [] \wedge *trv1'* \neq [] \wedge *trv2'* \neq [] \wedge
(*Suc* 0 < *length* *trv1* \vee *ww1* \leq *w1*) \wedge
(*Suc* 0 < *length* *trv1'* \vee *w1'* < *ww1*) \wedge
(*Suc* 0 < *length* *trv2* \vee *ww2* \leq *w2*) \wedge
(*Suc* 0 < *length* *trv2'* \vee *w2'* < *ww2*)
using φ φ 1 **unfolding** φ -def φ a-def **by** *auto*

show ?thesis
apply(*rule* *exI*[*of* - *w1*']) **apply**(*rule* *exI*[*of* - *w2*'])
apply(*rule* *exI*[*of* - *butlast* *trv1* @ *trv1'*]) **apply**(*rule* *exI*[*of* - *butlast* *trv2* @ *trv2'*])
apply(*rule* *exI*[*of* - *statA*']) **apply**(*rule* *exI*[*of* - *statO*'])
unfolding φ a-def **apply**(*intro* *conjI*)

```

subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def by auto
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def by auto
subgoal using len
by simp (metis Suc-lessI add-is-1 diff-is-0-eq length-greater-0-conv linorder-not-less
order-trans trans-less-add2)
subgoal using len
by simp (metis Suc-leI le-add-diff-inverse2 length-greater-0-conv nless-le order-le-less-trans trans-less-add2)
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def ssv1
using Van.validFromS-append by auto
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def ssv2
using Van.validFromS-append by auto
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def Van.S.map-filter Opt.S.map-filter

apply(subst tr1l) apply(subst butlast-append) by simp
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def Van.S.map-filter Opt.S.map-filter

apply(subst tr2l) apply(subst butlast-append) by simp
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def Van.A.map-filter Opt.A.map-filter

apply(subst trv1l) apply(subst trv2l)
apply(subst butlast-append) apply simp apply(subst butlast-append) by simp
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def apply simp
apply(cases Opt.O tr1 = Opt.O tr2, simp-all) apply clarify
using status.exhaust by (metis (full-types)+)
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def apply simp
apply(cases Opt.O tr1 = Opt.O tr2, simp-all) apply clarify
apply (smt (verit, del-insts) status.exhaust)
by (metis Opt.O.eq-Nil-iff nev(1) nev(2))
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def by simp
subgoal using  $\varphi$   $\varphi 1$  unfolding  $\varphi$ -def  $\varphi a$ -def by simp
subgoal using  $\varphi 1$  trv12'NE tr14NE unfolding  $\varphi$ -def  $\varphi a$ -def lastt-def by simp

```

qed

```

lemma unwindCond-ex- $\varphi a$ -aux2:
assumes unwind: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta$  w w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: (statA = Diff  $\longrightarrow$  statO = Diff)
and v3': Opt.validFromS s1 (tr1 @ [s1',s1'']) and v4': Opt.validFromS s2 (tr2 @ [s2',s2''])
and i: isIntO s1' isIntO s2'
and A34: getActO s1' = getActO s2'
and nev: never isIntO tr1 never isIntO tr2
shows  $\exists w1' w2' trv1 trv2 statAA statOO.$ 
   $\varphi a$   $\Delta$   $\infty$  w1 w2 w1' w2' statA s1 (tr1 @ [s1',s1'']) s2 (tr2 @ [s2',s2'']) statAA
  statO sv1 trv1 sv2 trv2 statOO

```

proof–

have 0 : $\text{lastt } s1 \ (tr1 \ \#\# \ s1') = s1' \ \text{lastt } s2 \ (tr2 \ \#\# \ s2') = s2'$
unfolding *lastt-def* **by** *auto*
show *?thesis*
apply(*rule unwindCond-ex- φ a'-aux[OF unwind Δ r stat, of tr1 $\#\#$ s1' tr2 $\#\#$ s2', unfolded 0, simplified]*)
using *assms* **by** *auto*
qed

lemma *lastt-snoc[simp]*: $\text{lastt } s1 \ (tr1 \ @ \ [s1'']) = s1''$
unfolding *lastt-def* **by** *auto*

lemma *lastt-snoc2[simp]*: $\text{lastt } s1 \ (tr1 \ @ \ [s1', s1'']) = s1''$
unfolding *lastt-def* **by** *auto*

lemma *append-snoc2*: $tr1 \ @ \ [s1', s1''] = (tr1 \ \#\# \ s1') \ \#\# \ s1''$
by *auto*

definition $\varphi' \ \Delta \ w1 \ w2 \ w1' \ w2' \ \text{statA } s1 \ tr1 \ s1' \ s1'' \ s2 \ tr2 \ s2' \ s2'' \ \text{statAA } \text{statO}$
 $sv1 \ trv1 \ sv1'' \ sv2 \ trv2 \ sv2'' \ \text{statOO} \equiv$
 $(trv1 \neq [] \vee w1' < w1) \wedge (trv2 \neq [] \vee w2' < w2) \wedge$
 $\text{Van.validFromS } sv1 \ (trv1 \ \#\# \ sv1'') \wedge \text{Van.validFromS } sv2 \ (trv2 \ \#\# \ sv2'') \wedge$
 $\text{Van.S } (trv1 \ \#\# \ sv1'') = \text{Opt.S } ((tr1 \ \#\# \ s1') \ \#\# \ s1'') \wedge \text{Van.S } (trv2 \ \#\# \ sv2'')$
 $= \text{Opt.S } ((tr2 \ \#\# \ s2') \ \#\# \ s2'') \wedge$
 $\text{Van.A } (trv1 \ \#\# \ sv1'') = \text{Van.A } (trv2 \ \#\# \ sv2'') \wedge$
 $(\text{statO} = \text{Eq} \longrightarrow (\text{statOO} = \text{Diff}) = (\text{Van.O } (trv1 \ \#\# \ sv1'') \neq \text{Van.O } (trv2 \ \#\# \ sv2''))) \wedge$
 $(\text{statA} = \text{Eq} \longrightarrow (\text{statAA} = \text{Diff}) = (\text{Opt.O } ((tr1 \ \#\# \ s1') \ \#\# \ s1'') \neq \text{Opt.O } ((tr2 \ \#\# \ s2') \ \#\# \ s2''))) \wedge$
 $(\text{statO} = \text{Diff} \longrightarrow \text{statOO} = \text{Diff}) \wedge (\text{statAA} = \text{Diff} \longrightarrow \text{statOO} = \text{Diff}) \wedge$
 $\Delta \ \infty \ w1' \ w2' \ s1'' \ s2'' \ \text{statAA } sv1'' \ sv2'' \ \text{statOO}$

proposition *unwindCond-ex- φ'* :

assumes *unwind*: *unwindCond Δ and Δ : $\Delta \ w \ w1 \ w2 \ s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$*
and *r*: *reachO s1 reachO s2 reachV sv1 reachV sv2*
and *stat*: *statA = Diff \longrightarrow statO = Diff*
and *v3'*: *Opt.validFromS s1 ((tr1 $\#\#$ s1') $\#\#$ s1'')* **and** *v4'*: *Opt.validFromS s2 ((tr2 $\#\#$ s2') $\#\#$ s2'')*
and *i*: *isIntO s1' isIntO s2'*
and *A34*: *getActO s1' = getActO s2'*
and *nev*: *never isIntO tr1 never isIntO tr2*
shows $\exists w1' \ w2' \ trv1 \ sv1'' \ trv2 \ sv2'' \ \text{statAA } \text{statOO}$.
 $\varphi' \ \Delta \ w1 \ w2 \ w1' \ w2' \ \text{statA } s1 \ tr1 \ s1' \ s1'' \ s2 \ tr2 \ s2' \ s2'' \ \text{statAA } \text{statO } sv1 \ trv1$
 $sv1'' \ sv2 \ trv2 \ sv2'' \ \text{statOO}$
using *unwindCond-ex- φ a-aux2[unfolding φ -def, unfolded lastt-snoc lastt-snoc2 append-snoc2, OF assms]*
unfolding *φ a-def* **apply**(*elim exE*) **subgoal for** $w1' \ w2' \ trv1 \ trv2 \ \text{statAA } \text{statOO}$
apply(*cases trv1 rule: rev-cases*)

subgoal by auto
apply(cases trv2 rule: rev-cases)
subgoal by auto
subgoal unfolding φ' -def **apply simp by blast . .**

definition $\chi^3 \Delta w (w1::enat) w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2$
 $statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv2 > Suc\ 0 \vee w2' \leq w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $never\ isSecV\ (butlast\ trv1) \wedge$
 $isSecV\ (lastt\ sv1\ trv1) \wedge getSecV\ (lastt\ sv1\ trv1) = getSecO\ (lastt\ s1\ tr1) \wedge$
 $never\ isSecV\ (butlast\ trv2) \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ (lastt\ s1\ tr1)\ s2\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma χ^3 -final:

assumes unw: unwindCond Δ
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and vtr1: Opt.validFromS s1 tr1
and χ^3 : $\chi^3 \Delta w w1 w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2 statOO$
shows (finalV (lastt sv1 trv1) \longleftrightarrow finalO (lastt s1 tr1)) \wedge (finalV (lastt sv2 trv2)
 \longleftrightarrow finalO s2)

proof –

have rsv12: Van.validFromS sv1 trv1 \longrightarrow reachV (lastt sv1 trv1)
 $Van.validFromS\ sv2\ trv2 \longrightarrow reachV\ (lastt\ sv2\ trv2)$ **using** r
by (simp add: Van.reach-validFromS-reach lastt-def)+
have rs1: Opt.validFromS s1 tr1 \longrightarrow reachO (lastt s1 tr1)
using r
by (simp add: Opt.reach-validFromS-reach lastt-def)+
show ?thesis **using** χ^3 [unfolded χ^3 -def] rsv12 rs1 **using** unw[unfolded unwind-
Cond-def, rule-format,
of lastt s1 tr1 s2 lastt sv1 trv1 lastt sv2 trv2 w w1' w2' statAA statOO]
using vtr1 <reachO s2> **by auto**
qed

lemma χ^3 -completedFrom: unwindCond $\Delta \implies$

reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies
Opt.validFromS s1 tr1 \implies completedFromO s1 tr1 \implies
 $\chi^3 \Delta w w1 w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2 statOO$
 \implies completedFromV sv1 trv1 \wedge completedFromV sv2 trv2
by (metis Van.final-not-isSec χ^3 -def χ^3 -final completedFromO-lastt)

lemma unwindCond-ex- χ^3 :

assumes unwind: unwindCond Δ
and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and r: reachO s1 reachO s2 reachV sv1 reachV sv2

and $vtr1$: *Opt.validFromS* $s1$ $tr1$
and $nis1$: $\neg isIntO$ $s1$ **and** $nis2$: $\neg isIntO$ $s2$
and $inter3$: *never isIntO* $tr1$
and sec : *never isSecO* (*butlast* $tr1$) *isSecO* (*lastt* $s1$ $tr1$)
shows $\exists w' w1' w2' trv1 trv2 statOO. \chi^3 \Delta w' w1 w2 w1' w2' s1 tr1 s2 statA sv1$
 $trv1 sv2 trv2 statOO$
using *assms*(2-)
proof(*induction length* $tr1$ w
arbitrary: $w1 w2 s1 s2 statA sv1 sv2 statO tr1$ *rule*: *less2-induct'*)
case (*less* $w tr1 w1 w2 s1 s2 statA sv1 sv2 statO$)
note $vtr1 = less(8)$

note $\Delta = \langle \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \rangle$
note $nis1 = less(9)$ **note** $nis2 = less(10)$
note $inter3 = less(11)$
note $sec3 = less(12,13)$
note $r34 = less.premis(2,3)$ **note** $r12 = less.premis(4,5)$
note $r = r34 r12$
note $r3 = r34(1)$ **note** $r4 = r34(2)$ **note** $r1 = r12(1)$ **note** $r2 = r12(2)$

have $i34$: $statA = Eq \longrightarrow isIntO s1 = isIntO s2$
and $f34$: $finalO s1 = finalO s2 \wedge finalV sv1 = finalO s1 \wedge finalV sv2 = finalO$
 $s2$
using Δ *unwind*[*unfolded unwindCond-def*] r **by** *auto*

have *proact-match*: ($\exists v < w. proact \Delta v w1 w2 s1 s2 statA sv1 sv2 statO$) \vee *react*
 $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$
using Δ *unwind*[*unfolded unwindCond-def*] r **by** *auto*
show *?case using proact-match proof safe*
fix v **assume** $v: v < w$
assume *proact* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
thus *?thesis unfolding proact-def proof safe*
assume $sv1$: $\neg isSecV sv1 \neg isIntV sv1$ **and** *move-1* $\Delta v w1 w2 s1 s2 statA$
 $sv1 sv2 statO$
then obtain $sv1'$
where 0 :*validTransV* ($sv1, sv1'$)
and Δ : $\Delta v w1 w2 s1 s2 statA sv1' sv2 statO$
unfolding *move-1-def* **by** *auto*
have $r1'$: *reachV* $sv1'$ **using** $r1$ 0 **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w' w1' w2' trv1 trv2 statOO$ **where** χ^3 : $\chi^3 \Delta w' w1 w2 w1' w2' s1$
 $tr1 s2 statA sv1' trv1 sv2 trv2 statOO$
using *less(2)*[*OF* v , *of* $tr1 w1 w2 s1 s2 statA sv1' sv2 statO$,
simplified, *OF* $\Delta r34 r1' r2 vtr1 nis1 nis2 inter3 sec3$] **by** *auto*
show *?thesis apply*(*rule exI*[*of - w'*]) **apply**(*rule exI*[*of - w1'*]) **apply**(*rule*
 exI [*of - w2'*]) **apply**(*rule exI*[*of - sv1 # trv1*]) **apply**(*rule exI*[*of - trv2*])
using χ^3 0 $sv1$ **unfolding** χ^3 -*def* **by** *auto*
next
assume $sv2$: $\neg isSecV sv2 \neg isIntV sv2$ **and** *move-2* $\Delta v w1 w2 s1 s2 statA$
 $sv1 sv2 statO$


```

then obtain sv2'
where 0: validTransV (sv2,sv2')
and Δ: Δ v w1 w2 s1 s2 statA sv1 sv2' statO
unfolding move-2-def by auto
have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w' w1' w2' trv1 trv2 statOO where χ3: χ3 Δ w' w1 w2 w1' w2' s1
tr1 s2 statA sv1 trv1 sv2' trv2 statOO
using less(2)[OF v, of tr1 w1 w2 s1 s2 statA sv1 sv2' statO,
simplified, OF Δ r34 r1 r2' vtr1 nis1 nis2 inter3 sec3] by auto
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using χ3 0 sv2 unfolding χ3-def by auto
next
assume sv12: ¬ isSecV sv1 ¬ isSecV sv2 Van.eqAct sv1 sv2
and move-12 Δ v w1 w2 s1 s2 statA sv1 sv2 statO
then obtain sv1' sv2' statO'
where 0: statO' = sstatO' statO sv1 sv2
validTransV (sv1,sv1') ¬ isSecV sv1
validTransV (sv2,sv2') ¬ isSecV sv2
Van.eqAct sv1 sv2
and Δ: Δ v w1 w2 s1 s2 statA sv1' sv2' statO'
unfolding move-12-def by auto
have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
fst-conv snd-conv)+

obtain w' w1' w2' trv1 trv2 statOO where χ3: χ3 Δ w' w1 w2 w1' w2' s1
tr1 s2 statA sv1' trv1 sv2' trv2 statOO
using less(2)[OF v, of tr1 w1 w2 s1 s2 statA sv1' sv2' statO',
simplified, OF Δ r34 r12' vtr1 nis1 nis2 inter3 sec3] by auto
show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
apply(rule exI[of - statOO])
using χ3 0 sv12 unfolding χ3-def sstatO'-def
by (auto simp: Van.eqAct-def)
qed
next
assume m: react Δ w1 w2 s1 s2 statA sv1 sv2 statO
define statA' where statA': statA' = sstatA' statA s1 s2
show ?thesis
proof(cases length tr1 ≤ Suc 0)
case True
hence tr1e: tr1 = [] ∨ tr1 = [s1]
by (metis Opt.validFromS-singl-iff Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
vtr1)
hence Opt.A tr1 = [] by (simp add: True)
have is1: isSecO s1
by (metis last.simps lastt-def sec3(2) tr1e)
hence ¬ finalO s1 using Opt.final-not-isSec by blast
then obtain s1' where s13': validTransO (s1, s1') unfolding Opt.final-def

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by auto
  hence isv1: isSecV sv1  $\wedge$  getSecV sv1 = getSecO s1 using m is1 nis1
  unfolding react-def match1-def eqSec-def by auto
  show ?thesis using tr1e isv1 apply-
    apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule exI[of - w2])
    apply(rule exI[of - [sv1]], rule exI[of - [sv2]], rule exI[of - statO])
    using tr1e
    using f34  $\Delta$  by (clarsimp simp:  $\chi$ 3-def lastt-def)
next
case False
then obtain s13 tr1' where tr1: tr1 = s13 # tr1' and tr1'NE: tr1'  $\neq$  []
  by (cases tr1, auto)
  have s13[simp]: s13 = s1 using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
    by (simp add: Opt.validFromS-Cons-iff tr1)
  obtain s1' where
    trn3: validTransO (s1, s1') and
    tr1': Opt.validFromS s1' tr1' using  $\langle$ Opt.validFromS s1 tr1 $\rangle$ 
  unfolding tr1 s13 by (metis tr1'NE Simple-Transition-System.validFromS-Cons-iff)
  have r3': reachO s1' using r3 trn3 by (metis Opt.reach.Step fst-conv snd-conv)
  have f3:  $\neg$  finalO s1 using Opt.final-def trn3 by blast
  hence f4:  $\neg$  finalO s2 using f34 by blast
  have nev3': never isIntO tr1'
  using inter3 tr1 tr1'NE by auto
  have isAO3:  $\neg$  isIntO s1 by (simp add: nis1)
  have O33': Opt.O tr1 = Opt.O tr1' Opt.A tr1 = Opt.A tr1'
  using isAO3 unfolding tr1 by auto
  have m: match1  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
  have  $(\exists w1' < w1. \exists w2' < w2. \neg isSecO s1 \wedge \Delta \infty w1' w2' s1' s2 statA sv1$ 
sv2 statO)  $\vee$ 
     $(\exists w2' < w2. eqSec sv1 s1 \wedge \neg isIntV sv1 \wedge match1-1 \Delta \infty w2' s1 s1' s2$ 
statA sv1 sv2 statO)  $\vee$ 
     $(eqSec sv1 s1 \wedge \neg isSecV sv2 \wedge Van.eqAct sv1 sv2 \wedge match1-12 \Delta \infty \infty$ 
s1 s1' s2 statA sv1 sv2 statO)
  using m isAO3 trn3 unfolding match1-def by auto
  thus ?thesis
proof safe
  fix w1'' w2'' assume w12': w1'' < w1 w2'' < w2
  assume  $\neg isSecO s1$  and  $\Delta$ :  $\Delta \infty w1'' w2'' s1' s2 statA sv1 sv2 statO$ 
  hence S3: Opt.S tr1' = Opt.S tr1 unfolding tr1 by auto
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi$ 3:  $\chi$ 3  $\Delta$  w' w1'' w2'' w1' w2'
s1' tr1' s2 statA sv1 trv1 sv2 trv2 statOO
  using less(1)[of tr1', OF -  $\Delta$  r3' r4 r12 -] unfolding tr1
  by simp (metis Opt.S.eq-Nil-iff(2) S3 Opt.validFromS-def  $\langle$  $\neg isSecO s1$  $\rangle$ 
last.simps
    lastt-def list-all-hd nev3' nis2 s13 sec3(1) sec3(2) tr1 tr1')
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
  using  $\chi$ 3 O33' unfolding  $\chi$ 3-def tr1 Van.completedFrom-def

```

```

using Van.validFromS-Cons  $tr1'NE$   $tr1'$   $trn3$  isAO3  $w12'$  by auto
next
fix  $w2''$  assume  $w2'$ :  $w2'' < w2$ 
assume  $trn13$ : eqSec  $sv1$   $s1$  and
 $Atrn1$ :  $\neg isIntV$   $sv1$  and  $match1-1$   $\Delta \infty w2'' s1 s1' s2$  statA  $sv1$   $sv2$  statO
then obtain  $sv1'$  where
 $trn1$ : validTransV ( $sv1, sv1'$ ) and
 $\Delta$ :  $\Delta \infty \infty w2'' s1' s2$  statA  $sv1'$   $sv2$  statO
unfolding  $match1-1-def$  by auto
have  $r1'$ : reachV  $sv1'$  using  $r1$   $trn1$  by (metis Van.reach.Step fst-conv snd-conv)
obtain  $w' w1' w2' trv1 trv2$  statOO where  $\chi3$ :  $\chi3 \Delta w' \infty w2'' w1' w2'$ 
 $s1' tr1' s2$  statA  $sv1' trv1 sv2 trv2$  statOO

using less(1)[of  $tr1'$ ,  $OF - \Delta r3' r4' r1' r2$ , unfolded O33', simplified]
using less.premis  $tr1' f3 f4 tr1'NE trn3 O33'(1)$ 
unfolding  $tr1$ 
by simp (metis Opt.validFromS-def list-all-hd)
show ?thesis apply(rule exI[of -  $w'$ ]) apply(rule exI[of -  $w1'$ ]) apply(rule
exI[of -  $w2'$ ]) apply(rule exI[of -  $sv1 \# trv1$ ]) apply(rule exI[of -  $trv2$ ])
using  $\chi3$  O33' unfolding  $\chi3-def$   $tr1$  Van.completedFrom-def
using Van.validFromS-Cons  $trn1$   $tr1'NE$   $tr1'$   $trn3$ 
using isAO3  $Atrn1$  eqSec-S-Cons  $trn13$   $w2'$ 
by simp (metis Opt.S.Nil-iff Opt.S.eq-Nil-iff(1) eqSec-def nless-le order-le-less-trans
 $s13$  sec3(1) tr1)
next
assume  $sv2$ :  $\neg isSecV$   $sv2$  and  $trn13$ : eqSec  $sv1$   $s1$  and
 $Atrn12$ : Van.eqAct  $sv1$   $sv2$  and  $match1-12$   $\Delta \infty \infty s1 s1' s2$  statA  $sv1$   $sv2$ 
statO
then obtain  $sv1' sv2' statO'$  where
 $statO'$ :  $statO' = sstatO' statO$   $sv1$   $sv2$  and
 $trn1$ : validTransV ( $sv1, sv1'$ ) and
 $trn2$ : validTransV ( $sv2, sv2'$ ) and
 $\Delta$ :  $\Delta \infty \infty \infty s1' s2$  statA  $sv1' sv2' statO'$ 
unfolding  $match1-12-def$  by auto
have  $r12'$ : reachV  $sv1'$  reachV  $sv2'$ 
using  $r1$   $trn1$   $r2$   $trn2$  by (metis Van.reach.Step fst-conv snd-conv)+
obtain  $w' w1' w2' trv1 trv2$  statOO where  $\chi3$ :  $\chi3 \Delta w' \infty \infty w1' w2' s1'$ 
 $tr1' s2$  statA  $sv1' trv1 sv2' trv2$  statOO
using less(1)[of  $tr1'$ ,  $OF - \Delta r3' r4' r12'$ , unfolded O33', simplified]
using less.premis  $tr1' f3 f4 tr1'NE trn3 O33'(1)$  unfolding  $tr1$  statO'
sstatO'-def
by simp (metis Simple-Transition-System.validFromS-def list-all-hd)+
show ?thesis apply(rule exI[of -  $w'$ ]) apply(rule exI[of -  $w1'$ ]) apply(rule
exI[of -  $w2'$ ]) apply(rule exI[of -  $sv1 \# trv1$ ]) apply(rule exI[of -  $sv2 \# trv2$ ])
using  $\chi3$  O33'  $tr1'NE$   $sv2$ 
using Van.validFromS-Cons  $trn1$   $trn2$ 
using isAO3  $Atrn12$  eqSec-S-Cons  $trn13$   $f3$   $f34$   $s13$   $tr1'$   $trn3$ 
unfolding  $\chi3-def$   $tr1$  Van.completedFrom-def Van.eqAct-def

```

```

    using Van.A.Cons-unfold eqSec-def sec3(1) tr1 by auto
  qed
  qed
  qed
  qed

```

definition $\chi3a$ **where** $\chi3a \Delta w (w1::enat) w2 w1' w2' s1 s1' s2 statAA sv1 trv1 sv2 trv2 statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv2 > Suc\ 0 \vee w2' < w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = [getSecO\ s1] \wedge$
 $never\ isSecV\ (butlast\ trv2) \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ s1'\ s2\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma *unwindCond-ex- $\chi3a$ -getSec:*

assumes *unwind:* *unwindCond* Δ

and $\Delta: \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and $r34: reachO\ s1\ reachO\ s2$ **and** $r12: reachV\ sv1\ reachV\ sv2$

and $v: validTransO\ (s1, s1')$

and $ii3: \neg isIntO\ s1$

and $is1: isSecO\ s1$ **and** $isv13: isSecV\ sv1\ getSecO\ s1 = getSecV\ sv1$

shows $\exists w1'\ w2'\ trv1\ trv2\ statOO.$

$\chi3a \Delta \infty w1\ w2\ w1'\ w2'\ s1\ s1'\ s2\ statA\ sv1\ trv1\ sv2\ trv2\ statOO$

using $\Delta\ r12\ isv13$

proof(*induction* w *arbitrary:* $w1\ w2\ sv1\ sv2\ statO$ *rule:* *less-induct*)

case (*less* $w\ w1\ w2\ sv1\ sv2\ statO$)

note $\Delta = \langle \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \rangle$

note $r12 = less.premis(2,3)$

note $r1 = r12(1)$ **note** $r2 = r12(2)$

note $r = r34\ r12$

note $isv13 = \langle isSecV\ sv1 \rangle \langle getSecO\ s1 = getSecV\ sv1 \rangle$

have $f34: finalO\ s1 = finalO\ s2 \wedge finalV\ sv1 = finalO\ s1 \wedge finalV\ sv2 = finalO\ s2$

using $\Delta\ unwind[unfolded\ unwindCond-def]$ r **by** *auto*

have *proact-match:* $(\exists v < w. proact\ \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO) \vee react\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

using $\Delta\ unwind[unfolded\ unwindCond-def]$ r **by** *auto*

show *?case* **using** *proact-match* **proof** *safe*

fix v **assume** $v: v < w$

assume *proact* $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

thus *?thesis* **unfolding** *proact-def* **proof** *safe*

assume $sv1: \neg isSecV\ sv1 \neg isIntV\ sv1$ **and** *move-1* $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

hence *False* **using** *isv13* **by** *blast*

thus *?thesis* **by** *auto*

next

```

assume sv2:  $\neg$  isSecV sv2  $\neg$  isIntV sv2 and move-2  $\Delta$  v w1 w2 s1 s2 statA
sv1 sv2 statO
then obtain sv2'
where 0: validTransV (sv2,sv2')
and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2' statO
unfolding move-2-def by auto
have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
obtain w1' w2' trv1 trv2 statOO where
 $\chi$ 3a:  $\chi$ 3a  $\Delta$   $\infty$  w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv2' trv2 statOO
using less(1)[OF v  $\Delta$  r1 r2' isv13] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
using  $\chi$ 3a 0 sv2 unfolding  $\chi$ 3a-def by auto
next
assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
and move-12  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2 statO
hence False using isv13 by blast
thus ?thesis by auto
qed
next
assume m: react  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
have m: match1  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
have ( $\exists$  w1' w2'. w1' < w1  $\wedge$  w2' < w2  $\wedge$   $\neg$  isSecO s1  $\wedge$   $\Delta$   $\infty$  w1' w2' s1' s2
statA sv1 sv2 statO)  $\vee$ 
( $\exists$  w2' < w2. eqSec sv1 s1  $\wedge$   $\neg$  isIntV sv1  $\wedge$  match1-1  $\Delta$   $\infty$  w2' s1 s1' s2
statA sv1 sv2 statO)  $\vee$ 
(eqSec sv1 s1  $\wedge$   $\neg$  isSecV sv2  $\wedge$  Van.eqAct sv1 sv2  $\wedge$  match1-12  $\Delta$   $\infty$   $\infty$ 
s1 s1' s2 statA sv1 sv2 statO)
using m v ii3 unfolding match1-def by auto

thus ?thesis
apply(elim disjE exE)
subgoal for w1' w2' using is1 by auto
subgoal for w2' apply(rule exI[of -  $\infty$ ]) apply(rule exI[of - w2'])
unfolding match1-1-def apply(elim conjE exE) subgoal for sv1'
apply(rule exI[of - [sv1,sv1']]) apply(rule exI[of - [sv2]])
apply(rule exI[of - statO])
using is1 isv13 unfolding  $\chi$ 3a-def
by (auto simp : sstatA'-def lastt-def) .
subgoal apply(rule exI[of -  $\infty$ ]) apply(rule exI[of -  $\infty$ ])
unfolding match1-12-def apply(elim conjE exE) subgoal for sv1' sv2'
apply(rule exI[of - [sv1,sv1']]) apply(rule exI[of - [sv2,sv2']])
apply(rule exI[of - sstatO' statO sv1 sv2])
using is1 isv13 unfolding  $\chi$ 3a-def
by (auto simp : sstatA'-def sstatO'-def lastt-def Van.eqAct-def) . .
qed
qed

```

definition $\chi^3b \Delta w (w1::enat) w2 w1' w2' s1 tr1 s2 statAA sv1 trv1 sv2 trv2$
 $statOO \equiv$
 $trv1 \neq [] \wedge$
 $trv2 \neq [] \wedge (length\ trv2 > Suc\ 0 \vee w2' < w2) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge$
 $Van.validFromS\ sv2\ trv2 \wedge$
 $Van.S\ trv1 = Opt.S\ tr1 \wedge$
 $never\ isSecV\ (butlast\ trv2) \wedge Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ (lastt\ s1\ tr1)\ s2\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma *unwindCond-ex- χ^3b -aux:*

assumes *unwind:* $unwindCond\ \Delta$

and $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ **and**

r: $reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$

and *tr1NE:* $tr1 \neq []$

and *v3':* $Opt.validFromS\ s1\ (tr1\ \#\#\ s1')$

and *nis1:* $\neg\ isIntO\ s1$ **and** *nis2:* $\neg\ isIntO\ s2$

and *ninter3':* $never\ isIntO\ (tr1\ \#\#\ s1')$

and *sec:* $never\ isSecO\ (butlast\ tr1)\ isSecO\ (lastt\ s1\ tr1)$

shows $\exists\ w1'\ w2'\ trv1\ trv2\ statOO.\ \chi^3b\ \Delta\ \infty\ w1\ w2\ w1'\ w2'\ s1\ (tr1\ \#\#\ s1')\ s2$
 $statA\ sv1\ trv1\ sv2\ trv2\ statOO$

proof–

have *v3:* $Opt.validFromS\ s1\ tr1$ **and** *s13':* $validTransO\ (lastt\ s1\ tr1, s1')$

apply (*metis* *v3'* $Opt.validFromS-def\ Opt.validS-append1\ Nil-is-append-conv\ hd-append2$)

by (*metis* $Opt.validFromS-def\ Opt.validS-validTrans\ lastt-def\ list.sel(1)\ not-Cons-self2$
 $snoc-eq-iff-butlast\ tr1NE\ v3'$)

have *ninter3:* $never\ isIntO\ tr1$ **and** *nis1':* $\neg\ isIntO\ s1'$

using *ninter3'* **by** *auto*

obtain *ww* *ww1* *ww2* *trv1* *trv2* *statOO* **where** $\chi^3:$ $\chi^3\ \Delta\ ww\ w1\ w2\ ww1\ ww2\ s1$
 $tr1\ s2\ statA\ sv1\ trv1\ sv2\ trv2\ statOO$

using $unwindCond-ex-\chi^3[OF\ unwind\ \Delta\ r\ v3\ nis1\ nis2\ ninter3\ sec]$ **by** *auto*

have *trv12NE:* $trv1 \neq []\ trv2 \neq []$ **using** χ^3 **unfolding** χ^3-def **by** *auto*

define *ss1* *ssv1* *ssv2* **where** *ss1:* $ss1 \equiv lastt\ s1\ tr1$

and *ssv1:* $ssv1 \equiv lastt\ sv1\ trv1$ **and** *ssv2:* $ssv2 \equiv lastt\ sv2\ trv2$

have *ss1l:* $ss1 = last\ tr1$ **by** (*simp* *add:* $lastt-def\ ss1\ tr1NE$)

have *tr1l:* $tr1 = butlast\ tr1\ @\ [ss1]$ **by** (*simp* *add:* $ss1\ tr1NE$)

have *ssv1l:* $ssv1 = last\ trv1$ **using** χ^3 **unfolding** χ^3-def **by** (*metis* $lastt-def$
 $ssv1$)

have *trv1l:* $trv1 = butlast\ trv1\ @\ [ssv1]$ **by** (*simp* *add:* $ssv1\ trv12NE(1)$)

have *ssv2l:* $ssv2 = last\ trv2$ **using** χ^3 **unfolding** χ^3-def **by** (*metis* $lastt-def$
 $ssv2$)

have *trv2l:* $trv2 = butlast\ trv2\ @\ [ssv2]$ **by** (*simp* *add:* $ssv2l\ trv12NE(2)$)

have $iss1[simp]: isSecO\ ss1$ **using** $sec(2)$ **unfolding** $ss1$ **by** $auto$
have $issv1[simp]: isSecV\ ssv1$ **and** $gissv13[simp]: getSecO\ ss1 = getSecV\ ssv1$
using χ^3 **unfolding** $\chi^3\text{-def}\ ssv1\ ss1$ **by** $auto$

have $niss1: \neg isIntO\ ss1$
by ($metis\ list\text{-all}\text{-append}\ list\text{-all}\text{-simps}(1)\ ninter3\ tr1l$)

have $rss1: reachO\ ss1$ **and** $rssv12: reachV\ ssv1\ reachV\ ssv2$
using $Opt.reach\text{-validFromS}\text{-reach}\ r\ ss1\ tr1NE\ v3$ **apply** $blast$
apply ($metis\ Van.reach\text{-validFromS}\text{-reach}\ \chi^3\text{-def}\ \chi^3\ r(3)\ ssv1l$)
by ($metis\ Van.reach\text{-validFromS}\text{-reach}\ \chi^3\text{-def}\ \chi^3\ r(4)\ ssv2l$)

have $\Delta: \Delta\ ww\ ww1\ ww2\ ss1\ s2\ statA\ ssv1\ ssv2\ statOO$
using χ^3 **unfolding** $\chi^3\text{-def}\ ss1[symmetric]\ ssv1[symmetric]\ ssv2[symmetric]$ **by**
 $auto$

have $s13': validTransO\ (ss1, s1')$
by ($simp\ add: s13'\ ss1$)

note $vs13 = s13'[unfolded\ ss1[symmetric]]$
obtain $w1'\ w2'\ trv1'\ trv2'\ statO'$ **where**
 $\chi^3a: \chi^3a\ \Delta\ \infty\ ww1\ ww2\ w1'\ w2'\ ss1\ s1'\ s2\ statA\ ssv1\ trv1'\ ssv2\ trv2'\ statO'$
using $unwindCond\text{-ex}\ \chi^3a\text{-getSec}[OF\ unwind\ \Delta\ rss1\ r(2)\ rssv12\ s13'\ niss1\ iss1\ issv1\ gissv13]$
by $blast$

have $trv12'NE: trv1' \neq []\ trv2' \neq []$ **using** χ^3a **unfolding** $\chi^3a\text{-def}$ **by** $auto$

have $[simp]: Van.O\ (butlast\ trv1\ @\ trv1') = Van.O\ trv1\ @\ Van.O\ trv1'$
using $trv12'NE$ **unfolding** $\chi^3\text{-def}\ Van.O.\text{map}\text{-filter}\ Opt.O.\text{map}\text{-filter}$ **apply** ($subst\ butlast\text{-append}$) **by** $simp$

have $[simp]: Van.O\ (butlast\ trv2\ @\ trv2') = Van.O\ trv2\ @\ Van.O\ trv2'$
using $trv12'NE$ **unfolding** $\chi^3\text{-def}\ Van.O.\text{map}\text{-filter}\ Opt.O.\text{map}\text{-filter}$ **apply** ($subst\ butlast\text{-append}$) **by** $simp$

have $Van.A\ trv1' = Van.A\ trv2'$ **using** χ^3a **unfolding** $\chi^3a\text{-def}$ **by** $auto$
moreover **have** $length\ (Van.O\ trv1') = length\ (Van.A\ trv1') \wedge length\ (Van.O\ trv2') = length\ (Van.A\ trv2')$
unfolding $Van.A.\text{map}\text{-filter}\ Van.O.\text{map}\text{-filter}$ **by** $auto$
ultimately **have** $length\ (Van.O\ trv1') = length\ (Van.O\ trv2')$ **by** $auto$
hence $[simp]: Van.O\ trv1\ @\ Van.O\ trv1' = Van.O\ trv2\ @\ Van.O\ trv2' \longleftrightarrow$
 $Van.O\ trv1 = Van.O\ trv2 \wedge Van.O\ trv1' = Van.O\ trv2'$ **by** $auto$

show $?thesis$
apply ($rule\ exI[of\ -\ w1']$) **apply** ($rule\ exI[of\ -\ w2']$)
apply ($rule\ exI[of\ -\ butlast\ trv1\ @\ trv1']$) **apply** ($rule\ exI[of\ -\ butlast\ trv2\ @\ trv2']$)
apply ($rule\ exI[of\ -\ statO']$)

unfolding $\chi^3 b\text{-def}$ **apply**(*intro conjI*)
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$ **by** *auto*
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$ **by** *auto*
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$
by *simp (metis Simple-Transition-System.fromS-eq-Nil Simple-Transition-System.toS-fromS-nonSingl Van.toS-Nil diff-add-inverse2 linorder-not-less order-le-less-trans trans-less-add2 zero-less-diff)*

subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$ *ssv1*
using *Van.validFromS-append* **by** *auto*
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$ *ssv2*
using *Van.validFromS-append* **by** *auto*
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$ **unfolding** *Van.S.map-filter*
Opt.S.map-filter
apply(*subst tr1l*) **apply**(*subst butlast-append*)
by *simp (metis Opt.S.map-filter Opt.S.eq-Nil-iff(2) Van.S.map-filter Van.S.eq-Nil-iff(2) sec(1))*
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$
by (*simp add: butlast-append*)
subgoal using $\chi^3 \chi^3 a$ **unfolding** $\chi^3\text{-def} \chi^3 a\text{-def}$ *Van.A.map-filter Opt.A.map-filter*

apply(*subst trv1l*) **apply**(*subst trv2l*) **by** (*simp add: butlast-append*)
subgoal using $\chi^3 a$ *trv12'NE tr1NE* **unfolding** $\chi^3 a\text{-def}$ *lastt-def* **by** *simp* .
qed

lemma *unwindCond-ex- $\chi^3 b\text{-aux2}$* :
assumes *unwind: unwindCond Δ*
and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$
and *r*: *reachO s1 reachO s2 reachV sv1 reachV sv2*
and *v3'*: *Opt.validFromS s1 (tr1 @ [s1',s1'])*
and *nis1*: $\neg isIntO s1$ **and** *nis2*: $\neg isIntO s2$
and *ninter3'*: *never isIntO (tr1 @ [s1',s1'])*
and *sec*: *never isSecO tr1 isSecO s1'*
shows $\exists w1' w2' trv1 trv2 statOO. \chi^3 b \Delta \infty w1 w2 w1' w2' s1 (tr1 @ [s1',s1'])$
 $s2 statA sv1 trv1 sv2 trv2 statOO$
proof –
have *0*: *lastt s1 (tr1 ## s1') = s1'*
unfolding *lastt-def* **by** *auto*
show *?thesis*
using *unwindCond-ex- $\chi^3 b\text{-aux}$ [OF unwind Δ r, of tr1 ## s1', unfolded 0, simplified]*
using *assms* **by** *auto*
qed

definition $\chi^3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statAA sv1 trv1 sv1'' sv2 trv2$
 $sv2'' statOO \equiv$
 $Van.validFromS sv1 (trv1 ## sv1'') \wedge Van.validFromS sv2 (trv2 ## sv2'') \wedge$
 $Van.S (trv1 ## sv1'') = Opt.S ((tr1 ## s1') ## s1'') \wedge never isSecV trv2 \wedge$

$Van.A (trv1 \#\# sv1'') = Van.A (trv2 \#\# sv2'') \wedge$
 $trv1 \neq [] \wedge (trv2 \neq [] \vee w2' < w2) \wedge$
 $\Delta \infty w1' w2' s1'' s2 statAA sv1'' sv2'' statOO$

proposition *unwindCond-ex- χ^3'* :

assumes *unwind*: *unwindCond* Δ

and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ **and**

r: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *v3'*: *Opt.validFromS* *s1* ((*tr1* $\#\#$ *s1'*) $\#\#$ *s1''*)

and *nis1*: $\neg isIntO$ *s1* **and** *nis2*: $\neg isIntO$ *s2*

and *ninter3'*: *never isIntO* ((*tr1* $\#\#$ *s1'*) $\#\#$ *s1''*)

and *sec*: *never isSecO* *tr1* *isSecO* *s1'*

shows $\exists w1' w2' trv1 sv1'' trv2 sv2'' statOO. \chi^3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$

using *unwindCond-ex- χ^3b -aux2*[*unfolded φ -def*, *unfolded lastt-snoc lastt-snoc2 append-snoc2*, *OF assms*]

unfolding *χ^3b -def* **apply**(*elim exE*) **subgoal for** *w1' w2' trv1 trv2 statOO*

apply(*cases trv1 rule: rev-cases*)

subgoal by *auto*

subgoal for *trv1' sv1''* **apply**(*cases trv2 rule: rev-cases*)

subgoal by *auto*

subgoal for *trv2' sv2''* **unfolding** *χ^3' -def*

apply(*rule exI[of - w1']*) **apply**(*rule exI[of - w2']*)

apply(*rule exI[of - trv1']*) **apply**(*rule exI[of - sv1'']*)

apply(*rule exI[of - trv2']*) **apply**(*rule exI[of - sv2'']*)

apply(*rule exI[of - statOO]*)

by *simp* (*metis Opt.S.Nil-iff Opt.S.eq-Nil-iff(1) Van.S.simps(4) append-snoc2 list-all-append sec(2)*)

self-append-conv2 snoc-eq-iff-butlast) . . .

definition $\omega^3 \Delta w1 w2 w1' w2' s1 s1' s2 statAA sv1 trv1 sv1' sv2 trv2 sv2'$
statOO \equiv

Van.validFromS *sv1* (*trv1* $\#\#$ *sv1'*) \wedge *Van.validFromS* *sv2* (*trv2* $\#\#$ *sv2'*) \wedge
never isSecV *trv1* \wedge *never isSecV* *trv2* \wedge

Van.A (*trv1* $\#\#$ *sv1'*) = *Van.A* (*trv2* $\#\#$ *sv2'*) \wedge

(*trv1* $\neq [] \vee w1' < w1$) \wedge (*trv2* $\neq [] \vee w2' < w2$) \wedge

$\Delta \infty w1' w2' s1' s2 statAA sv1' sv2' statOO$

proposition *unwindCond-ex- ω^3* :

assumes *unwind*: *unwindCond* Δ

and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$

and *r34*: *reachO* *s1* *reachO* *s2* **and** *r12*: *reachV* *sv1* *reachV* *sv2*

and *v3*: *validTransO* (*s1*, *s1'*)

and *nis1*: $\neg isIntO$ *s1* $\neg isIntO$ *s1'* $\neg isSecO$ *s1*

and *nis2*: $\neg isIntO$ *s2*

shows $\exists w1' w2' trv1 sv1' trv2 sv2' statOO. \omega^3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2' statOO$

```

using  $\Delta$   $r12$ 
proof(induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct)
  case (less w w1 w2 sv1 sv2 statO)
  note  $\Delta = \langle \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \rangle$ 
  note  $r12 = less.prem(2,3)$ 
  note  $r1 = r12(1)$  note  $r2 = r12(2)$ 
  note  $r = r34 r12$ 

  have  $f34: finalO s1 = finalO s2 \wedge finalV sv1 = finalO s1 \wedge finalV sv2 = finalO$ 
 $s2$ 
    using  $\Delta$  unwind[unfolded unwindCond-def]  $r$  by auto

  have proact-match: ( $\exists v < w. proact \Delta v w1 w2 s1 s2 statA sv1 sv2 statO$ )  $\vee react$ 
 $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$ 
    using  $\Delta$  unwind[unfolded unwindCond-def]  $r$  by auto
    show ?case using proact-match proof safe
    fix  $v$  assume  $v: v < w$ 
    assume proact  $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$ 
    thus ?thesis unfolding proact-def proof safe
    assume  $sv1: \neg isSecV sv1 \neg isIntV sv1$  and move-1  $\Delta v w1 w2 s1 s2 statA$ 
 $sv1 sv2 statO$ 
    then obtain  $sv1'$  where  $0: validTransV (sv1, sv1')$  and  $\Delta: \Delta v w1 w2 s1$ 
 $s2 statA sv1' sv2 statO$ 
    unfolding move-1-def by auto
    have  $r1': reachV sv1'$  using  $r1 0$  by (metis Van.reach.Step fst-conv snd-conv)
    obtain  $w1' w2' trv1 sv1'' trv2 sv2'' statOO$  where
     $\omega3: \omega3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1' trv1 sv1'' sv2 trv2 sv2'' statOO$ 

    using less(1)[OF v  $\Delta r1' r2$ ] by auto
    show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule
 $exI$ [of - sv1 # trv1]) apply(rule exI[of - sv1''])
    apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
    using  $\omega3 0 sv1$  unfolding  $\omega3$ -def by auto
  next
    assume  $sv2: \neg isSecV sv2 \neg isIntV sv2$  and move-2  $\Delta v w1 w2 s1 s2 statA$ 
 $sv1 sv2 statO$ 
    then obtain  $sv2'$ 
    where  $0: validTransV (sv2, sv2')$ 
    and  $\Delta: \Delta v w1 w2 s1 s2 statA sv1 sv2' statO$ 
    unfolding move-2-def by auto
    have  $r2': reachV sv2'$  using  $r2 0$  by (metis Van.reach.Step fst-conv snd-conv)
    obtain  $w1' w2' trv1 sv1' trv2 sv2'' statOO$  where
     $\omega3: \omega3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2' trv2 sv2'' statOO$ 

    using less(1)[OF v  $\Delta r1 r2'$ ] by auto
    show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
    apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
    using  $\omega3 0 sv2$  unfolding  $\omega3$ -def by auto

```

```

next
  assume  $sv1: \neg isSecV\ sv1$  and  $sv2: \neg isSecV\ sv2$  and
  move-12  $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$  and
  sv12:  $Van.eqAct\ sv1\ sv2$ 
  then obtain  $sv1'\ sv2'\ statO'$ 
  where  $statO': statO' = sstatO'\ statO\ sv1\ sv2$ 
  and  $0: validTransV\ (sv1,sv1')\ validTransV\ (sv2,sv2')$ 
  and  $\Delta: \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1'\ sv2'\ statO'$ 
  unfolding move-12-def by auto
  have  $r1': reachV\ sv1'$  and  $r2': reachV\ sv2'$  using  $r1\ r2\ 0$ 
  by (metis  $Van.reach.Step\ fst-conv\ snd-conv$ )
  obtain  $w1'\ w2'\ trv1\ sv1''\ trv2\ sv2''\ statOO$  where
   $\omega3: \omega3\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s1'\ s2\ statA\ sv1'\ trv1\ sv1''\ sv2'\ trv2\ sv2''\ statOO$ 

  using less(1)[ $OF\ v\ \Delta\ r1'\ r2'$ ] by auto
  show ?thesis apply(rule exI[ $of - w1'$ ]) apply(rule exI[ $of - w2'$ ])
  apply(rule exI[ $of - sv1\ \# trv1$ ]) apply(rule exI[ $of - sv1''$ ])
  apply(rule exI[ $of - sv2\ \# trv2$ ]) apply(rule exI[ $of - sv2''$ ])
  using  $\omega3\ 0\ sv1\ sv2\ sv12$  unfolding  $\omega3-def\ statO'$  by (auto simp:  $Van.eqAct-def$ )
qed
next
  assume  $m: react\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
  have  $m: match1\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$  using  $m$  unfolding
  react-def by auto
  have  $(\exists w1'\ w2'. w1' < w1 \wedge w2' < w2 \wedge \neg isSecO\ s1 \wedge \Delta\ \infty\ w1'\ w2'\ s1'\ s2\ statA\ sv1\ sv2\ statO) \vee$ 
   $(\exists w2' < w2. eqSec\ sv1\ s1 \wedge \neg isIntV\ sv1 \wedge match1-1\ \Delta\ \infty\ w2'\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO) \vee$ 
   $(eqSec\ sv1\ s1 \wedge \neg isSecV\ sv2 \wedge Van.eqAct\ sv1\ sv2 \wedge match1-12\ \Delta\ \infty\ \infty\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO)$ 
  using  $m\ v3\ nis1$  unfolding match1-def by auto

  thus ?thesis
  apply(elim disjE exE)
  subgoal for  $w1'\ w2'$ 
  apply(rule exI[ $of - w1'$ ]) apply(rule exI[ $of - w2'$ ])
  apply(rule exI[ $of - []$ ]) apply(rule exI[ $of - sv1$ ])
  apply(rule exI[ $of - []$ ]) apply(rule exI[ $of - sv2$ ])
  apply(rule exI[ $of - statO$ ]) unfolding  $\omega3-def$ 
  by auto
  subgoal for  $w2'$ 
  apply(rule exI[ $of - \infty$ ]) apply(rule exI[ $of - w2'$ ])
  unfolding match1-1-def apply(elim conjE exE) subgoal for  $sv1'$ 
  apply(rule exI[ $of - [sv1]$ ]) apply(rule exI[ $of - sv1'$ ])
  apply(rule exI[ $of - []$ ]) apply(rule exI[ $of - sv2$ ])
  apply(rule exI[ $of - statO$ ])
  unfolding  $\omega3-def$  using  $nis1(3)$  by (auto simp: eqSec-def) .
  subgoal
  apply(rule exI[ $of - \infty$ ]) apply(rule exI[ $of - \infty$ ])

```

unfolding *match1-12-def* **apply**(*elim conjE exE*) **subgoal for** *sv1' sv2'*
apply(*rule exI[of - [sv1]]*) **apply**(*rule exI[of - sv1']*)
apply(*rule exI[of - [sv2]]*) **apply**(*rule exI[of - sv2']*)
apply(*rule exI[of - sstatO' statO sv1 sv2]*)
unfolding *ω3-def* **using** *nis1(3)* **apply** (*auto simp: eqSec-def*
sstatA'-def sstatO'-def lastt-def Van.eqAct-def) . . .

qed
qed

definition $\chi_4 \Delta w w1 (w2::\text{enat}) w1' w2' s1 s2 tr2 \text{statAA } sv1 \text{trv1 } sv2 \text{trv2}$
 $\text{statOO} \equiv$
 $\text{trv1} \neq [] \wedge \text{trv2} \neq [] \wedge (\text{length } \text{trv1} > \text{Suc } 0 \vee w1' \leq w1) \wedge$
 $\text{Van.validFromS } sv1 \text{trv1} \wedge \text{Van.validFromS } sv2 \text{trv2} \wedge$
 $\text{never isSecV (butlast } \text{trv1}) \wedge$
 $\text{never isSecV (butlast } \text{trv2}) \wedge$
 $\text{isSecV (lastt } sv2 \text{trv2}) \wedge \text{getSecV (lastt } sv2 \text{trv2}) = \text{getSecO (lastt } s2 \text{tr2}) \wedge$
 $\text{Van.A } \text{trv1} = \text{Van.A } \text{trv2} \wedge$
 $\Delta w w1' w2' s1 (\text{lastt } s2 \text{tr2}) \text{statAA (lastt } sv1 \text{trv1) (lastt } sv2 \text{trv2) statOO}$

lemma χ_4 -*final*:

assumes *unw*: *unwindCond* Δ

and *r*: *reachO* *s1 reachO s2 reachV sv1 reachV sv2*

and *vtr2*: *Opt.validFromS s2 tr2*

and χ_4 : $\chi_4 \Delta w w1 w2 w1' w2' s1 s2 tr2 \text{statAA } sv1 \text{trv1 } sv2 \text{trv2 statOO}$

shows (*finalV (lastt sv1 trv1)* \longleftrightarrow *finalO s1*) \wedge (*finalV (lastt sv2 trv2)* \longleftrightarrow *finalO (lastt s2 tr2)*)

proof–

have *rsv12*: *Van.validFromS sv1 trv1* \longrightarrow *reachV (lastt sv1 trv1)*

Van.validFromS sv2 trv2 \longrightarrow *reachV (lastt sv2 trv2)* **using** *r*

by (*simp add: Van.reach-validFromS-reach lastt-def*) $+$

have *rs2*: *Opt.validFromS s2 tr2* \longrightarrow *reachO (lastt s2 tr2)*

using *r*

by (*simp add: Opt.reach-validFromS-reach lastt-def*) $+$

show *?thesis* **using** χ_4 [*unfolded* χ_4 -*def*] *rsv12 rs2* **using** *unw*[*unfolded* *unwindCond-def, rule-format,*

of s1 lastt s2 tr2 lastt sv1 trv1 lastt sv2 trv2 w w1' w2' statAA statOO]

using *vtr2* \langle *reachO s1* \rangle **by** *auto*

qed

lemma χ_4 -*completedFrom*: *unwindCond* $\Delta \implies$

reachO s1 \implies *reachO s2* \implies *reachV sv1* \implies *reachV sv2* \implies

Opt.validFromS s2 tr2 \implies *completedFromO s2 tr2* \implies

$\chi_4 \Delta w w1 w2 w1' w2' s1 s2 tr2 \text{statAA } sv1 \text{trv1 } sv2 \text{trv2 statOO}$

\implies *completedFromV sv1 trv1* \wedge *completedFromV sv2 trv2*

by (*metis* *Van.final-not-isSec* χ_4 -*def* χ_4 -*final* *completedFromO-lastt*)

proposition *unwindCond-ex- χ_4* :
assumes *unwind*: *unwindCond* Δ
and Δ : Δ *w w1 w2 s1 s2 statA sv1 sv2 statO*
and *r*: *reachO s1 reachO s2 reachV sv1 reachV sv2*
and *vtr2*: *Opt.validFromS s2 tr2*
and *nis2*: \neg *isIntO s1* **and** *nis2*: \neg *isIntO s2*
and *inter4*: *never isIntO tr2*
and *sec*: *never isSecO (butlast tr2) isSecO (lastt s2 tr2)*
shows $\exists w' w1' w2' trv1 trv2 statOO. \chi_4 \Delta w' w1 w2 w1' w2' s1 s2 tr2 statA sv1 trv1 sv2 trv2 statOO$
using *assms(2-)*
proof(*induction length tr2 w*
arbitrary: *w1 w2 s1 s2 statA sv1 sv2 statO tr2 rule: less2-induct'*)
case (*less w tr2 w1 w2 s1 s2 statA sv1 sv2 statO*)
note *vtr2 = less(8)*
note $\Delta = \langle \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \rangle$
note *nis1 = less(9)* **note** *nis2 = less(10)*
note *inter4 = less(11)*
note *sec4 = less(12,13)*
note *r34 = less.prem(2,3)* **note** *r12 = less.prem(4,5)*
note *r = r34 r12*
note *r3 = r34(1)* **note** *r4 = r34(2)* **note** *r1 = r12(1)* **note** *r2 = r12(2)*

have *i34*: *statA = Eq \longrightarrow isIntO s1 = isIntO s2*
and *f34*: *finalO s1 = finalO s2 \wedge finalV sv1 = finalO s1 \wedge finalV sv2 = finalO s2*
using Δ *unwind[unfolded unwindCond-def]* *r* **by** *auto*

have *proact-match*: ($\exists v < w. \text{proact } \Delta v w1 w2 s1 s2 statA sv1 sv2 statO$) \vee *react*
 $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$
using Δ *unwind[unfolded unwindCond-def]* *r* **by** *auto*
show *?case using proact-match proof safe*
fix *v* **assume** *v*: *v < w*
assume *proact* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
thus *?thesis unfolding proact-def proof safe*
assume *sv1*: \neg *isSecV sv1* \neg *isIntV sv1* **and** *move-1* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
then obtain *sv1'*
where *0*:*validTransV (sv1,sv1')*
and Δ : $\Delta v w1 w2 s1 s2 statA sv1' sv2 statO$
unfolding *move-1-def* **by** *auto*
have *r1'*: *reachV sv1'* **using** *r1 0* **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain *w' w1' w2' trv1 trv2 statOO* **where** χ_4 : $\chi_4 \Delta w' w1 w2 w1' w2' s1 s2 tr2 statA sv1' trv1 sv2 trv2 statOO$
using *less(2)[OF v, of tr2 w1 w2 s1 s2 statA sv1' sv2 statO,*
simplified, OF $\Delta r34 r1' r2 vtr2 nis1 nis2 inter4 sec4]$ **by** *auto*
show *?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - trv2])*
using $\chi_4 0 sv1$ **unfolding** χ_4 -*def* **by** *auto*

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next
  assume sv2:  $\neg$  isSecV sv2  $\neg$  isIntV sv2 and move-2  $\Delta$  v w1 w2 s1 s2 statA
  sv1 sv2 statO
  then obtain sv2'
  where 0: validTransV (sv2,sv2')
  and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2' statO
  unfolding move-2-def by auto
  have r2': reachV sv2' using r2 0 by (metis Van.reach.Step fst-conv snd-conv)
  obtain w1' w2' w' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4$   $\Delta$  w' w1 w2 w1' w2' s1
  s2 tr2 statA sv1 trv1 sv2' trv2 statOO
  using less(2)[OF v, of tr2 w1 w2 s1 s2 statA sv1 sv2' statO,
    simplified, OF  $\Delta$  r34 r1 r2' vtr2 nis1 nis2 inter4 sec4] by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
  exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\chi_4$  0 sv2 unfolding  $\chi_4$ -def by auto
  next
  assume sv12:  $\neg$  isSecV sv1  $\neg$  isSecV sv2 Van.eqAct sv1 sv2
  and move-12  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2 statO
  then obtain sv1' sv2' statO'
  where 0: statO' = sstatO' statO sv1 sv2
  validTransV (sv1,sv1')  $\neg$  isSecV sv1
  validTransV (sv2,sv2')  $\neg$  isSecV sv2
  Van.eqAct sv1 sv2
  and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1' sv2' statO'
  unfolding move-12-def by auto
  have r12': reachV sv1' reachV sv2' using r1 r2 0 by (metis Van.reach.Step
  fst-conv snd-conv)+
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4$   $\Delta$  w' w1 w2 w1' w2' s1
  s2 tr2 statA sv1' trv1 sv2' trv2 statOO
  using less(2)[OF v, of tr2 w1 w2 s1 s2 statA sv1' sv2' statO',
    simplified, OF  $\Delta$  r34 r12' vtr2 nis1 nis2 inter4 sec4] by auto
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
  exI[of - w2']) apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
  apply(rule exI[of - statOO])
  using  $\chi_4$  0 sv12 unfolding  $\chi_4$ -def sstatO'-def
  by (auto simp: Van.eqAct-def)
  qed
  next
  assume m: react  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
  define statA' where statA': statA' = sstatA' statA s1 s2
  show ?thesis
  proof(cases length tr2  $\leq$  Suc 0)
  case True
  hence tr2e: tr2 = []  $\vee$  tr2 = [s2]
  by (metis Opt.validFromS-def Suc-length-conv le-Suc-eq le-zero-eq length-0-conv
  list.sel(1) vtr2)
  hence Opt.A tr2 = [] by (simp add: True)
  have is2: isSecO s2
  by (metis last.simps lastt-def sec4(2) tr2e)

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hence  $\neg \text{finalO } s2$  using Opt.final-not-isSec by blast
then obtain  $s2'$  where  $s24'$ : validTransO ( $s2, s2'$ ) unfolding Opt.final-def
by auto
hence  $\text{isv2}: \text{isSecV } sv2 \wedge \text{getSecV } sv2 = \text{getSecO } s2$  using m is2 nis2
unfolding react-def match2-def eqSec-def by auto
show ?thesis using tr2e isv2 apply-
  apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule exI[of - w2])
  apply(rule exI[of - [sv1]], rule exI[of - [sv2]], rule exI[of - statO])
  using tr2e
  using f34  $\Delta$  by(clarsimp simp:  $\chi4$ -def lastt-def)
next
case False
then obtain  $s24$   $tr2'$  where  $tr2: tr2 = s24 \# tr2'$  and  $tr2'NE: tr2' \neq []$ 
  by (cases tr2, auto)
  have  $s24$ [simp]:  $s24 = s2$  using  $\langle \text{Opt.validFromS } s2 \text{ } tr2 \rangle$ 
    by (simp add: Opt.validFromS-Cons-iff tr2)
  obtain  $s2'$  where
     $trn4: \text{validTransO } (s2, s2')$  and
     $tr2': \text{Opt.validFromS } s2' \text{ } tr2'$  using  $\langle \text{Opt.validFromS } s2 \text{ } tr2 \rangle$ 
  unfolding tr2 s24 by (metis tr2'NE Simple-Transition-System.validFromS-Cons-iff)
  have  $r4'$ : reachO  $s2'$  using  $r4$   $trn4$  by (metis Opt.reach.Step fst-conv snd-conv)
  have  $f4$ :  $\neg \text{finalO } s2$  using Opt.final-def trn4 by blast
  hence  $f3$ :  $\neg \text{finalO } s1$  using f34 by blast
  have  $nev4'$ : never isIntO  $tr2'$ 
  using inter4 tr2 tr2'NE by auto
  have  $isAO4$ :  $\neg \text{isIntO } s2$  by (simp add: nis2)
  have  $O44'$ :  $\text{Opt.O } tr2 = \text{Opt.O } tr2' \text{ } \text{Opt.A } tr2 = \text{Opt.A } tr2'$ 
  using  $isAO4$  unfolding tr2 by auto
  have  $m$ : match2  $\Delta w1 w2 s1 s2 \text{ statA } sv1 sv2 \text{ statO}$  using m unfolding react-def by auto
  have  $(\exists w1' < w1. \exists w2' < w2. \neg \text{isSecO } s2 \wedge \Delta \infty w1' w2' s1 s2' \text{ statA } sv1 sv2 \text{ statO}) \vee$ 
     $(\exists w1' < w1. \text{eqSec } sv2 s2 \wedge \neg \text{isIntV } sv2 \wedge \text{match2-1 } \Delta w1' \infty s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO}) \vee$ 
     $(\text{eqSec } sv2 s2 \wedge \neg \text{isSecV } sv1 \wedge \text{Van.eqAct } sv1 sv2 \wedge \text{match2-12 } \Delta \infty \infty s1 s2 s2' \text{ statA } sv1 sv2 \text{ statO})$ 
  using m isAO4 trn4 unfolding match2-def by auto
  thus ?thesis
proof safe
  fix  $w1'' w2''$  assume  $w12'$ :  $w1'' < w1 w2'' < w2$ 
  assume  $\neg \text{isSecO } s2$  and  $\Delta: \Delta \infty w1'' w2'' s1 s2' \text{ statA } sv1 sv2 \text{ statO}$ 
  hence  $S4$ :  $\text{Opt.S } tr2' = \text{Opt.S } tr2$  unfolding tr2 by auto
  obtain  $w' w1' w2' \text{ trv1 } \text{trv2 } \text{statOO}$  where  $\chi4$ :  $\chi4 \Delta w' w1'' w2'' w1' w2'$ 
 $s1 s2' tr2' \text{ statA } sv1 \text{ trv1 } sv2 \text{ trv2 } \text{statOO}$ 
  using less(1)[of tr2', OF -  $\Delta r3 r4' r12$ ] unfolding tr2
  by simp (metis Opt.S.eq-Nil-iff(2) S4 Simple-Transition-System.validFromS-def last.simps lastt-def
  list.discI list-all-hd nev4' nis1 sec4(1) sec4(2) tr2 tr2' tr2'NE)
  show ?thesis apply(rule exI[of - w]) apply(rule exI[of - w1]) apply(rule

```

```

exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - trv2])
  using  $\chi_4$  O44' unfolding  $\chi_4$ -def tr2 Van.completedFrom-def
  using Van.validFromS-Cons tr2'NE tr2' trn4 isAO4 w12' by auto
next
  fix w1'' assume w1': w1'' < w1
  assume trn24: eqSec sv2 s2 and
  Atrn2:  $\neg$  isIntV sv2 and match2-1  $\Delta$  w1''  $\infty$  s1 s2 s2' statA sv1 sv2 statO
  then obtain sv2' where
  trn2: validTransV (sv2,sv2') and
   $\Delta$ :  $\Delta$   $\infty$  w1''  $\infty$  s1 s2' statA sv1 sv2' statO
  unfolding match2-1-def by auto
  have r2': reachV sv2' using r2 trn2 by (metis Van.reach.Step fst-conv
snd-conv)
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4$   $\Delta$  w' w1''  $\infty$  w1' w2'
s1 s2' tr2' statA sv1 trv1 sv2' trv2 statOO
  using less(1)[of tr2', OF -  $\Delta$  r3 r4' r1 r2', unfolded O44', simplified]
  using less.premis tr2' f3 f4 tr2'NE trn4 O44'(1)
  unfolding tr2
  by simp (metis Opt.validFromS-def list-all-hd)
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2']) apply(rule exI[of - trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\chi_4$  O44' unfolding  $\chi_4$ -def tr2 Van.completedFrom-def
  using Van.validFromS-Cons trn2 tr2'NE tr2' trn4
  using isAO4 Atrn2 eqSec-S-Cons trn24 w1'
  by simp (metis Opt.S.Nil-iff Opt.S.eq-Nil-iff(1) eqSec-def nless-le order-le-less-trans
s24 sec4(1) tr2)
next
  assume sv1:  $\neg$  isSecV sv1 and trn24: eqSec sv2 s2 and
  Atrn12: Van.eqAct sv1 sv2 and match2-12  $\Delta$   $\infty$   $\infty$  s1 s2 s2' statA sv1 sv2
statO
  then obtain sv1' sv2' statO' where
  statO': statO' = sstatO' statO sv1 sv2 and
  trn1: validTransV (sv1,sv1') and
  trn2: validTransV (sv2,sv2') and
   $\Delta$ :  $\Delta$   $\infty$   $\infty$   $\infty$  s1 s2' statA sv1' sv2' statO'
  unfolding match2-12-def by auto
  have r12': reachV sv1' reachV sv2'
  using r1 trn1 r2 trn2 by (metis Van.reach.Step fst-conv snd-conv)+
  obtain w' w1' w2' trv1 trv2 statOO where  $\chi_4$ :  $\chi_4$   $\Delta$  w'  $\infty$   $\infty$  w1' w2' s1
s2' tr2' statA sv1' trv1 sv2' trv2 statOO
  using less(1)[of tr2', OF -  $\Delta$  r3 r4' r12', unfolded O44', simplified]
  using less.premis tr2' f3 f4 tr2'NE trn4 O44'(1) unfolding tr2 statO'
sstatO'-def
  by simp (metis Simple-Transition-System.validFromS-def list-all-hd)
  show ?thesis apply(rule exI[of - w']) apply(rule exI[of - w1']) apply(rule
exI[of - w2'])
  apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv2 # trv2])
  using  $\chi_4$  O44' tr2'NE sv1
  using Van.validFromS-Cons trn1 trn2

```


using *isAO4 Atrn12 eqSec-S-Cons trn24 f3 f34 s24 tr2' trn4*
unfolding χ_4 -def tr2 *Van.complectedFrom-def Van.eqAct-def*
using *Van.A.Cons-unfold eqSec-def sec4(1) tr2* **by** *auto*
qed
qed
qed
qed

definition χ_4a **where** $\chi_4a \Delta w w1 (w2::enat) w1' w2' s1 s2 s2' statAA sv1 trv1$
 $sv2 trv2 statOO \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (length\ trv1 > Suc\ 0 \vee w1' < w1) \wedge$
 $Van.validFromS\ sv1\ trv1 \wedge Van.validFromS\ sv2\ trv2 \wedge$
 $never\ isSecV\ (butlast\ trv1) \wedge$
 $Van.S\ trv2 = [getSecO\ s2] \wedge$
 $Van.A\ trv1 = Van.A\ trv2 \wedge$
 $\Delta\ w\ w1'\ w2'\ s1\ s2'\ statAA\ (lastt\ sv1\ trv1)\ (lastt\ sv2\ trv2)\ statOO$

lemma *unwindCond-ex- χ_4a -getSec:*

assumes *unwind: unwindCond Δ*

and $\Delta: \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and $r34: reachO\ s1\ reachO\ s2$ **and** $r12: reachV\ sv1\ reachV\ sv2$

and $v: validTransO\ (s2, s2')$

and $ii4: \neg isIntO\ s2$

and $is2: isSecO\ s2$ **and** $isv24: isSecV\ sv2\ getSecO\ s2 = getSecV\ sv2$

shows $\exists w1' w2' trv1 trv2 statOO.$

$\chi_4a \Delta \infty w1\ w2\ w1' w2' s1\ s2\ s2' statA\ sv1\ trv1\ sv2\ trv2\ statOO$

using $\Delta\ r12\ isv24$

proof(*induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct*)

case (*less w w1 w2 sv1 sv2 statO*)

note $\Delta = \langle \Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \rangle$

note $r12 = less.premis(2,3)$

note $r1 = r12(1)$ **note** $r2 = r12(2)$

note $r = r34\ r12$

note $isv24 = \langle isSecV\ sv2 \rangle \langle getSecO\ s2 = getSecV\ sv2 \rangle$

have $f34: finalO\ s1 = finalO\ s2 \wedge finalV\ sv1 = finalO\ s1 \wedge finalV\ sv2 = finalO\ s2$

using $\Delta\ unwind[unfolding\ unwindCond-def]$ **r** **by** *auto*

have *proact-match: ($\exists v < w. proact\ \Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$) \vee react $\Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$*

using $\Delta\ unwind[unfolding\ unwindCond-def]$ **r** **by** *auto*

show *?case using proact-match proof safe*

fix v **assume** $v: v < w$

assume *proact $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$*

thus *?thesis unfolding proact-def proof safe*

assume $sv1: \neg isSecV\ sv1 \neg isIntV\ sv1$ **and** *move-1 $\Delta\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$*

then obtain $sv1'$

where 0 : *validTransV* ($sv1, sv1'$)
and Δ : $\Delta v w1 w2 s1 s2 statA sv1' sv2 statO$
unfolding *move-1-def* **by** *auto*
have $r1'$: *reachV* $sv1'$ **using** $r1 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w1' w2' trv1 trv2 statOO$ **where**
 $\chi4a$: $\chi4a \Delta \infty w1 w2 w1' w2' s1 s2 s2' statA sv1' trv1 sv2 trv2 statOO$
using *less(1)[OF v $\Delta r1' r2 isv24$]* **by** *auto*
show *?thesis* **apply**(*rule exI[of - $w1'$]*) **apply**(*rule exI[of - $w2'$]*) **apply**(*rule exI[of - $sv1 \# trv1$]*) **apply**(*rule exI[of - $trv2$]*)
using $\chi4a 0 sv1$ **unfolding** $\chi4a-def$ **by** *auto*
next
assume $sv2$: $\neg isSecV sv2 \neg isIntV sv2$ **and** *move-2* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
hence *False* **using** *isv24* **by** *blast*
thus *?thesis* **by** *auto*
next
assume $sv12$: $\neg isSecV sv1 \neg isSecV sv2 Van.eqAct sv1 sv2$
and *move-12* $\Delta v w1 w2 s1 s2 statA sv1 sv2 statO$
hence *False* **using** *isv24* **by** *blast*
thus *?thesis* **by** *auto*
qed
next
assume m : *react* $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$
have m : *match2* $\Delta w1 w2 s1 s2 statA sv1 sv2 statO$ **using** m **unfolding** *react-def* **by** *auto*
have $(\exists w1' w2'. w1' < w1 \wedge w2' < w2 \wedge \neg isSecO s2 \wedge \Delta \infty w1' w2' s1 s2' statA sv1 sv2 statO) \vee$
 $(\exists w1' < w1. eqSec sv2 s2 \wedge \neg isIntV sv2 \wedge match2-1 \Delta w1' \infty s1 s2 s2' statA sv1 sv2 statO) \vee$
 $(eqSec sv2 s2 \wedge \neg isSecV sv1 \wedge Van.eqAct sv1 sv2 \wedge match2-12 \Delta \infty \infty s1 s2 s2' statA sv1 sv2 statO)$
using $m v ii4$ **unfolding** *match2-def* **by** *auto*

thus *?thesis*
apply(*elim disjE exE*)
subgoal for $w1' w2'$ **using** *is2* **by** *auto*
subgoal for $w1'$ **apply**(*rule exI[of - $w1'$]*) **apply**(*rule exI[of - ∞]*)
unfolding *match2-1-def* **apply**(*elim conjE exE*) **subgoal for** $sv2'$
apply(*rule exI[of - $sv1$]*) **apply**(*rule exI[of - $sv2, sv2'$]*)
apply(*rule exI[of - $statO$]*)
using *is2 isv24* **unfolding** $\chi4a-def$
by (*auto simp : sstatA'-def lastt-def*) .
subgoal **apply**(*rule exI[of - ∞]*) **apply**(*rule exI[of - ∞]*)
unfolding *match2-12-def* **apply**(*elim conjE exE*) **subgoal for** $sv1' sv2'$
apply(*rule exI[of - $sv1, sv1'$]*) **apply**(*rule exI[of - $sv2, sv2'$]*)
apply(*rule exI[of - $sstatO' statO sv1 sv2$]*)
using *is2 isv24* **unfolding** $\chi4a-def$
by (*auto simp : sstatA'-def sstatO'-def lastt-def Van.eqAct-def*) . .
qed

qed

definition $\chi_4 b \Delta w w1 w2 w1' (w2'::\text{enat}) s1 s2 tr2 \text{statAA} sv1 trv1 sv2 trv2$
 $\text{statOO} \equiv$
 $trv1 \neq [] \wedge trv2 \neq [] \wedge (\text{length } trv1 > \text{Suc } 0 \vee w1' < w1) \wedge$
 $\text{Van.validFromS } sv1 trv1 \wedge \text{Van.validFromS } sv2 trv2 \wedge$
 $\text{never isSecV (butlast } trv1) \wedge$
 $\text{Van.S } trv2 = \text{Opt.S } tr2 \wedge$
 $\text{Van.A } trv1 = \text{Van.A } trv2 \wedge$
 $\Delta w w1' w2' s1 (\text{lastt } s2 tr2) \text{statAA} (\text{lastt } sv1 trv1) (\text{lastt } sv2 trv2) \text{statOO}$

lemma $\text{unwindCond-ex-}\chi_4 b\text{-aux}$:

assumes unwind : $\text{unwindCond } \Delta$

and Δ : $\Delta w w1 w2 s1 s2 \text{statA} sv1 sv2 \text{statO}$

and r : $\text{reachO } s1 \text{ reachO } s2 \text{ reachV } sv1 \text{ reachV } sv2$

and $tr2NE$: $tr2 \neq []$

and $v4'$: $\text{Opt.validFromS } s2 (tr2 \#\# s2')$

and $nis1$: $\neg \text{isIntO } s1$ **and** $nis2$: $\neg \text{isIntO } s2$

and $ninter4'$: $\text{never isIntO } (tr2 \#\# s2')$

and sec : $\text{never isSecO (butlast } tr2) \text{ isSecO (lastt } s2 tr2)$

shows $\exists w1' w2' trv1 trv2 \text{statOO}. \chi_4 b \Delta \infty w1 w2 w1' w2' s1 s2 (tr2 \#\# s2')$
 $\text{statA } sv1 trv1 sv2 trv2 \text{statOO}$

proof–

have $v4$: $\text{Opt.validFromS } s2 tr2$ **and** $s24'$: $\text{validTransO (lastt } s2 tr2, s2')$

apply ($\text{metis } v4' \text{Opt.validFromS-def Opt.validS-append1 Nil-is-append-conv hd-append2}$)

by ($\text{metis Opt.validFromS-def Opt.validS-validTrans append-is-Nil-conv lastt-def list.distinct(1) list.sel(1) tr2NE } v4'$)

have $ninter4$: $\text{never isIntO } tr2$ **and** $nis2'$: $\neg \text{isIntO } s2'$

using $ninter4'$ **by** auto

obtain $ww ww1 ww2 trv1 trv2 \text{statOO}$ **where** χ_4 : $\chi_4 \Delta ww w1 w2 ww1 ww2 s1$
 $s2 tr2 \text{statA} sv1 trv1 sv2 trv2 \text{statOO}$

using $\text{unwindCond-ex-}\chi_4$ [$OF \text{unwind } \Delta r v4 nis1 nis2 ninter4 sec$]

by auto

have $trv12NE$: $trv1 \neq [] \wedge trv2 \neq []$ **using** χ_4 **unfolding** $\chi_4\text{-def}$ **by** auto

define $ss2 ssv1 ssv2$ **where** $ss2$: $ss2 \equiv \text{lastt } s2 tr2$

and $ssv1$: $ssv1 \equiv \text{lastt } sv1 trv1$ **and** $ssv2$: $ssv2 \equiv \text{lastt } sv2 trv2$

have $ss2l$: $ss2 = \text{last } tr2$ **by** ($\text{simp add: lastt-def } ss2 tr2NE$)

have $tr2l$: $tr2 = \text{butlast } tr2 @ [ss2]$ **by** ($\text{simp add: } ss2l tr2NE$)

have $ssv1l$: $ssv1 = \text{last } trv1$ **using** χ_4 **unfolding** $\chi_4\text{-def}$ **by** ($\text{metis lastt-def } ssv1$)

have $trv1l$: $trv1 = \text{butlast } trv1 @ [ssv1]$ **by** ($\text{simp add: } ssv1l trv12NE(1)$)

have $ssv2l$: $ssv2 = \text{last } trv2$ **using** χ_4 **unfolding** $\chi_4\text{-def}$ **by** ($\text{metis lastt-def } ssv2$)

have $trv2l$: $trv2 = \text{butlast } trv2 @ [ssv2]$ **by** ($\text{simp add: } ssv2l trv12NE(2)$)

have $iss2[simp]: isSecO\ ss2$ **using** $sec(2)$ **unfolding** $ss2$ **by** *auto*
have $issv2[simp]: isSecV\ ssv2$ **and** $gissv24[simp]: getSecO\ ss2 = getSecV\ ssv2$
using χ_4 **unfolding** $\chi_4\text{-def}\ ssv2\ ss2$ **by** *auto*

have $niss2: \neg isIntO\ ss2$
by (*metis list-all-append list-all-simps(1) ninter4 tr2l*)

have $rss2: reachO\ ss2$ **and** $rssv12: reachV\ ssv1\ reachV\ ssv2$
using *Opt.reach-validFromS-reach* $r\ ss2l\ tr2NE\ v_4$ **apply** *blast*
unfolding $ssv1\ ssv2$ **using** $r(3,4)$ **using** χ_4 **unfolding** $\chi_4\text{-def}$
using *Van.reach-validFromS-reach* $ssv1\ ssv2\ ssv1l\ ssv2l$ **by** *auto metis+*

have $\Delta: \Delta\ ww\ ww1\ ww2\ s1\ ss2\ statA\ ssv1\ ssv2\ statOO$
using χ_4 **unfolding** $\chi_4\text{-def}\ ss2[symmetric]\ ssv1[symmetric]\ ssv2[symmetric]$ **by**
auto

have $s13': validTransO\ (ss2, s2')$
by (*simp add: s24' ss2*)

note $vs24 = s24'[unfolded\ ss2[symmetric]]$
obtain $w1'\ w2'\ trv1'\ trv2'\ statO'$ **where**
 $\chi_4a: \chi_4a\ \Delta\ \infty\ ww1\ ww2\ w1'\ w2'\ s1\ ss2\ s2'\ statA\ ssv1\ trv1'\ ssv2\ trv2'\ statO'$
using *unwindCond-ex- χ_4a -getSec[OF unwind $\Delta\ r(1)\ rss2\ rssv12\ s13'\ niss2\ iss2\ issv2\ gissv24$]*
by *blast*

have $trv12'NE: trv1' \neq []\ trv2' \neq []$ **using** χ_4a **unfolding** $\chi_4a\text{-def}$ **by** *auto*

have $[simp]: Van.O\ (butlast\ trv1\ @\ trv1') = Van.O\ trv1\ @\ Van.O\ trv1'$
using $trv12'NE$ **unfolding** $\chi_4\text{-def}\ Van.O.map-filter\ Opt.O.map-filter$ **apply** (*subst butlast-append*) **by** *simp*

have $[simp]: Van.O\ (butlast\ trv2\ @\ trv2') = Van.O\ trv2\ @\ Van.O\ trv2'$
using $trv12'NE$ **unfolding** $\chi_4\text{-def}\ Van.O.map-filter\ Opt.O.map-filter$ **apply** (*subst butlast-append*) **by** *simp*

have $Van.A\ trv1' = Van.A\ trv2'$ **using** χ_4a **unfolding** $\chi_4a\text{-def}$ **by** *auto*
moreover **have** $length\ (Van.O\ trv1') = length\ (Van.A\ trv1') \wedge length\ (Van.O\ trv2') = length\ (Van.A\ trv2')$
unfolding *Van.A.map-filter Van.O.map-filter* **by** *auto*
ultimately **have** $length\ (Van.O\ trv1') = length\ (Van.O\ trv2')$ **by** *auto*
hence $[simp]: Van.O\ trv1\ @\ Van.O\ trv1' = Van.O\ trv2\ @\ Van.O\ trv2' \longleftrightarrow$
 $Van.O\ trv1 = Van.O\ trv2 \wedge Van.O\ trv1' = Van.O\ trv2'$ **by** *auto*

show *?thesis*
apply (*rule exI[of - w1']*) **apply** (*rule exI[of - w2']*)
apply (*rule exI[of - butlast trv1 @ trv1']*) **apply** (*rule exI[of - butlast trv2 @ trv2']*)
apply (*rule exI[of - statO']*)

unfolding $\chi_4 b\text{-def}$ **apply**(*intro conjI*)
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$ **by** *auto*
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$ **by** *auto*
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$
by *simp (metis Simple-Transition-System.fromS-eq-Nil Van.toS-Nil Van.toS-fromS-nonSingl*
diff-add-inverse2 linorder-not-less order-le-less-trans trans-less-add2 zero-less-diff)

subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$ *ssv1*
using *Van.validFromS-append* **by** *auto*
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$ *ssv2*
using *Van.validFromS-append* **by** *auto*
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$
by (*simp add: butlast-append*)
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$ **unfolding** *Van.S.map-filter*
Opt.S.map-filter
apply(*subst tr2l*) **apply**(*subst butlast-append*)
by *simp (metis Opt.S.map-filter Opt.S.eq-Nil-iff Van.S.map-filter Van.S.eq-Nil-iff*
sec(1))
subgoal using $\chi_4 \chi_4 a$ **unfolding** $\chi_4\text{-def} \chi_4 a\text{-def}$ *Van.A.map-filter Opt.A.map-filter*

apply(*subst trv1l*) **apply**(*subst trv2l*)
apply(*subst butlast-append*) **apply** *simp* **apply**(*subst butlast-append*) **by** *simp*
subgoal using $\chi_4 a$ *trv12'NE tr2NE* **unfolding** $\chi_4 a\text{-def}$ *lastt-def* **by** *simp* .
qed

lemma *unwindCond-ex- $\chi_4 b\text{-aux2}$* :
assumes *unwind: unwindCond Δ*
and Δ : $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ **and**
r: reachO s1 reachO s2 reachV sv1 reachV sv2
and v_4' : *Opt.validFromS s2 (tr2 @ [s2',s2'])*
and $nis1$: $\neg isIntO s1$ **and** $nis2$: $\neg isIntO s2$
and $ninter_4'$: *never isIntO (tr2 @ [s2',s2'])*
and sec : *never isSecO tr2 isSecO s2'*
shows $\exists w1' w2' trv1 trv2 statOO. \chi_4 b \Delta \infty w1 w2 w1' w2' s1 s2 (tr2 @ [s2',s2'])$
statA sv1 trv1 sv2 trv2 statOO
proof –
have 0 : *lastt s2 (tr2 ## s2') = s2'*
unfolding *lastt-def* **by** *auto*
show *?thesis*
using *unwindCond-ex- $\chi_4 b\text{-aux}$ [OF unwind Δ r, of tr2 ## s2', unfolded 0, sim-*
plified]
using *assms* **by** *auto*
qed

definition $\chi_4' \Delta w1 w2 w1' (w2'::enat) s1 s2 tr2 s2' s2'' statAA sv1 trv1 sv1''$
 $sv2 trv2 sv2'' statOO \equiv$

$Van.validFromS\ sv1\ (trv1\ \#\#\ sv1'') \wedge Van.validFromS\ sv2\ (trv2\ \#\#\ sv2'') \wedge$
 $never\ isSecV\ (butlast\ (trv1\ \#\#\ sv1'')) \wedge$
 $Van.S\ (trv2\ \#\#\ sv2'') = Opt.S\ ((tr2\ \#\#\ s2')\ \#\#\ s2'') \wedge$
 $Van.A\ (trv1\ \#\#\ sv1'') = Van.A\ (trv2\ \#\#\ sv2'') \wedge$
 $trv2 \neq [] \wedge (trv1 \neq [] \vee w1' < w1) \wedge$
 $\Delta \infty w1'\ w2'\ s1\ s2''\ statAA\ sv1''\ sv2''\ statOO$

proposition *unwindCond-ex- χ_4'* :

assumes *unwind*: *unwindCond* Δ

and Δ : $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ **and**

r: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and v_4' : *Opt.validFromS* *s2* $((tr2\ \#\#\ s2')\ \#\#\ s2'')$

and *nis1*: $\neg\ isIntO\ s1$ **and** *nis2*: $\neg\ isIntO\ s2$

and *ninter* $_4'$: *never isIntO* $((tr2\ \#\#\ s2')\ \#\#\ s2'')$

and *sec*: *never isSecO* *tr2* *isSecO* *s2'*

shows $\exists\ w1'\ w2'\ trv1\ sv1''\ trv2\ sv2''\ statOO.\ \chi_4'\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s2\ tr2\ s2'\ s2''\ statA\ sv1\ trv1\ sv1''\ sv2\ trv2\ sv2''\ statOO$

using *unwindCond-ex- χ_4b -aux2*[*unfolded φ -def*, *unfolded lastt-snoc lastt-snoc2 append-snoc2*, *OF assms*]

unfolding χ_4b -*def* **apply**(*elim exE*) **subgoal for** $w1'\ w2'\ trv1\ trv2\ statOO$

apply(*cases trv1 rule: rev-cases*)

subgoal by *auto*

subgoal for $trv1'\ sv1''$ **apply**(*cases trv2 rule: rev-cases*)

subgoal by *auto*

subgoal for $trv2'\ sv2''$ **unfolding** χ_4' -*def*

apply(*rule exI*[*of - w1*]) **apply**(*rule exI*[*of - w2*])

apply(*rule exI*[*of - trv1*]) **apply**(*rule exI*[*of - sv1''*])

apply(*rule exI*[*of - trv2*]) **apply**(*rule exI*[*of - sv2''*])

apply(*rule exI*[*of - statOO*])

by *simp* (*metis Opt.S.Nil-iff Opt.S.eq-Nil-iff*(1) *Van.S.simps*(4) *butlast-append*

list.discI list-all-append sec(2) *self-append-conv2*) . . .

definition $\omega_4\ \Delta\ w1\ w2\ w1'\ (w2'::enat)\ s1\ s2\ s2'\ statAA\ sv1\ trv1\ sv1'\ sv2\ trv2\ sv2'\ statOO \equiv$

$Van.validFromS\ sv1\ (trv1\ \#\#\ sv1') \wedge Van.validFromS\ sv2\ (trv2\ \#\#\ sv2') \wedge$

$never\ isSecV\ trv1 \wedge never\ isSecV\ trv2 \wedge$

$Van.A\ (trv1\ \#\#\ sv1') = Van.A\ (trv2\ \#\#\ sv2') \wedge$

$(trv1 \neq [] \vee w1' < w1) \wedge (trv2 \neq [] \vee w2' < w2) \wedge$

$\Delta \infty w1'\ w2'\ s1\ s2'\ statAA\ sv1'\ sv2'\ statOO$

proposition *unwindCond-ex- ω_4* :

assumes *unwind*: *unwindCond* Δ

and Δ : $\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and r_{34} : *reachO* *s1* *reachO* *s2* **and** r_{12} : *reachV* *sv1* *reachV* *sv2*

and *nis1*: $\neg\ isIntO\ s1$

and v_4 : *validTransO* (s_2, s_2')
and nis_2 : $\neg isIntO\ s_2 \neg isIntO\ s_2' \neg isSecO\ s_2$
shows $\exists w_1' w_2' trv_1 sv_1' trv_2 sv_2' statOO. \omega_4 \Delta w_1 w_2 w_1' w_2' s_1 s_2 s_2' statA$
 $sv_1 trv_1 sv_1' sv_2 trv_2 sv_2' statOO$
using $\Delta\ r12$
proof(*induction w arbitrary: w1 w2 sv1 sv2 statO rule: less-induct*)
case (*less w w1 w2 sv1 sv2 statO*)
note $\Delta = \langle \Delta\ w\ w_1\ w_2\ s_1\ s_2\ statA\ sv_1\ sv_2\ statO \rangle$
note $r12 = less.prem_2(2,3)$
note $r1 = r12(1)$ **note** $r2 = r12(2)$
note $r = r34\ r12$

have $f34$: $finalO\ s_1 = finalO\ s_2 \wedge finalV\ sv_1 = finalO\ s_1 \wedge finalV\ sv_2 = finalO$
 s_2
using $\Delta\ unwind[unfolded\ unwindCond-def]\ r$ **by** *auto*

have *proact-match*: $(\exists v < w. proact\ \Delta\ v\ w_1\ w_2\ s_1\ s_2\ statA\ sv_1\ sv_2\ statO) \vee react$
 $\Delta\ w_1\ w_2\ s_1\ s_2\ statA\ sv_1\ sv_2\ statO$
using $\Delta\ unwind[unfolded\ unwindCond-def]\ r$ **by** *auto*
show *?case using proact-match proof safe*
fix v **assume** $v: v < w$
assume *proact* $\Delta\ v\ w_1\ w_2\ s_1\ s_2\ statA\ sv_1\ sv_2\ statO$
thus *?thesis unfolding proact-def proof safe*
assume $sv_1: \neg isSecV\ sv_1 \neg isIntV\ sv_1$ **and** *move-1* $\Delta\ v\ w_1\ w_2\ s_1\ s_2\ statA$
 $sv_1\ sv_2\ statO$
then obtain sv_1' **where** 0 : *validTransV* (sv_1, sv_1') **and** Δ : $\Delta\ v\ w_1\ w_2\ s_1$
 $s_2\ statA\ sv_1'\ sv_2\ statO$
unfolding *move-1-def* **by** *auto*
have $r1'$: *reachV* sv_1' **using** $r1\ 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w_1' w_2' trv_1 sv_1'' trv_2 sv_2' statOO$ **where**
 ω_4 : $\omega_4\ \Delta\ w_1\ w_2\ w_1' w_2' s_1\ s_2\ s_2' statA\ sv_1' trv_1 sv_1'' sv_2 trv_2 sv_2' statOO$

using *less(1)[OF v Δ r1' r2]* **by** *auto*
show *?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule*
 $exI[of - sv_1 \# trv_1]) apply(rule exI[of - sv_1''])$
 $apply(rule exI[of - trv_2]) apply(rule exI[of - sv_2'])$
using $\omega_4\ 0\ sv_1$ **unfolding** ω_4-def **by** *auto*
next
assume $sv_2: \neg isSecV\ sv_2 \neg isIntV\ sv_2$ **and** *move-2* $\Delta\ v\ w_1\ w_2\ s_1\ s_2\ statA$
 $sv_1\ sv_2\ statO$
then obtain sv_2'
where 0 : *validTransV* (sv_2, sv_2')
and Δ : $\Delta\ v\ w_1\ w_2\ s_1\ s_2\ statA\ sv_1\ sv_2' statO$
unfolding *move-2-def* **by** *auto*
have $r2'$: *reachV* sv_2' **using** $r2\ 0$ **by** (*metis Van.reach.Step fst-conv snd-conv*)
obtain $w_1' w_2' trv_1 sv_1' trv_2 sv_2'' statOO$ **where**
 ω_4 : $\omega_4\ \Delta\ w_1\ w_2\ w_1' w_2' s_1\ s_2\ s_2' statA\ sv_1 trv_1 sv_1' sv_2' trv_2 sv_2'' statOO$

using *less(1)[OF v Δ r1 r2']* **by** *auto*

```

show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
using  $\omega_4$  0 sv2 unfolding  $\omega_4$ -def by auto
next
assume sv1:  $\neg$  isSecV sv1 and sv2:  $\neg$  isSecV sv2 and
move-12  $\Delta$  v w1 w2 s1 s2 statA sv1 sv2 statO and
sv12: Van.eqAct sv1 sv2
then obtain sv1' sv2' statO'
where statO': statO' = sstatO' statO sv1 sv2
and 0: validTransV (sv1,sv1') validTransV (sv2,sv2'')
and  $\Delta$ :  $\Delta$  v w1 w2 s1 s2 statA sv1' sv2' statO'
unfolding move-12-def by auto
have r1': reachV sv1' and r2': reachV sv2' using r1 r2 0
by (metis Van.reach.Step fst-conv snd-conv)+
obtain w1' w2' trv1 sv1'' trv2 sv2'' statOO where
 $\omega_4$ :  $\omega_4$   $\Delta$  w1 w2 w1' w2' s1 s2 s2' statA sv1' trv1 sv1'' sv2' trv2 sv2'' statOO

using less(1)[OF v  $\Delta$  r1' r2'] by auto
show ?thesis apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - sv1 # trv1]) apply(rule exI[of - sv1''])
apply(rule exI[of - sv2 # trv2]) apply(rule exI[of - sv2''])
using  $\omega_4$  0 sv1 sv2 sv12 unfolding  $\omega_4$ -def statO' by (auto simp: Van.eqAct-def)
qed
next
assume m: react  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO
have m: match2  $\Delta$  w1 w2 s1 s2 statA sv1 sv2 statO using m unfolding
react-def by auto
have ( $\exists$  w1' w2'. w1' < w1  $\wedge$  w2' < w2  $\wedge$   $\neg$  isSecO s2  $\wedge$   $\Delta$   $\infty$  w1' w2' s1 s2'
statA sv1 sv2 statO)  $\vee$ 
( $\exists$  w1' < w1. eqSec sv2 s2  $\wedge$   $\neg$  isIntV sv2  $\wedge$  match2-1  $\Delta$  w1'  $\infty$  s1 s2 s2'
statA sv1 sv2 statO)  $\vee$ 
 $\neg$  isSecV sv1  $\wedge$  eqSec sv2 s2  $\wedge$  Van.eqAct sv1 sv2  $\wedge$  match2-12  $\Delta$   $\infty$   $\infty$ 
s1 s2 s2' statA sv1 sv2 statO
using m v4 nis2 unfolding match2-def by auto

thus ?thesis
apply(elim disjE exE)
subgoal for w1' w2' apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - []]) apply(rule exI[of - sv1])
apply(rule exI[of - []]) apply(rule exI[of - sv2])
apply(rule exI[of - statO]) unfolding  $\omega_4$ -def
by auto
subgoal for w1' apply(rule exI[of - w1']) apply(rule exI[of -  $\infty$ ])
unfolding match2-1-def apply(elim conjE exE) subgoal for sv2'
apply(rule exI[of - []]) apply(rule exI[of - sv1])
apply(rule exI[of - [sv2]]) apply(rule exI[of - sv2'])
apply(rule exI[of - statO])
unfolding  $\omega_4$ -def using nis2(3) by (auto simp: eqSec-def) .

```



```

subgoal apply(rule exI[of -  $\infty$ ]) apply(rule exI[of -  $\infty$ ])
unfolding match2-12-def apply(elim conjE exE) subgoal for sv1' sv2'
apply(rule exI[of - [sv1]]) apply(rule exI[of - sv1 ^])
apply(rule exI[of - [sv2]]) apply(rule exI[of - sv2 ^])
apply(rule exI[of - sstatO' statO sv1 sv2])
unfolding  $\omega_4$ -def using nis2(3) apply (auto simp: eqSec-def
sstatA'-def sstatO'-def lastt-def Van.eqAct-def) . . .
qed
qed

```

```

definition  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'  $\equiv$ 
  ltr1 = lappend (llist-of (tr1 ## s1')) (s1'' $ ltr1')  $\wedge$ 
  ltr2 = lappend (llist-of (tr2 ## s2')) (s2'' $ ltr2')  $\wedge$ 
  Opt.validFromS s1 ((tr1 ## s1') ## s1'')  $\wedge$  Opt.validFromS s2 ((tr2 ## s2')
## s2'')  $\wedge$ 
  never isIntO tr1  $\wedge$  never isIntO tr2  $\wedge$ 
  isIntO s1'  $\wedge$  isIntO s2'  $\wedge$  getActO s1' = getActO s2'  $\wedge$ 
  Opt.lvalidFromS s1'' (s1'' $ ltr1')  $\wedge$  Opt.lcompletedFrom s1'' (s1'' $ ltr1')  $\wedge$ 
  Opt.lvalidFromS s2'' (s2'' $ ltr2')  $\wedge$  Opt.lcompletedFrom s2'' (s2'' $ ltr2')  $\wedge$ 
  Opt.lA (s1'' $ ltr1') = Opt.lA (s2'' $ ltr2')

```

lemma isIntO- $\varphi\varphi$:

```

assumes vltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and vltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A: Opt.lA ltr1 = Opt.lA ltr2 and inter3:  $\neg$  lnever isIntO ltr1
shows  $\exists$  tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'.  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2
s2' s2'' ltr2'

```

proof–

```

have 03:  $\exists$  s $\in$ lset ltr1. isIntO s using inter3 unfolding llist.pred-set by auto
define ttr1 where ttr1: ttr1  $\equiv$  ltakeUntil isIntO ltr1
define lltr1' where lltr1': lltr1'  $\equiv$  ldropUntil isIntO ltr1
have ltr1: ltr1 = lappend (llist-of ttr1) lltr1'
unfolding ttr1 lltr1' lappend-ltakeUntil-ldropUntil[OF 03] ..
have 13: ttr1  $\neq$  []  $\wedge$  never isIntO (butlast ttr1)  $\wedge$  isIntO (last ttr1)
unfolding ttr1
using ltakeUntil-last[OF 03] ltakeUntil-not-Nil[OF 03] ltakeUntil-never-butlast[OF
03] by simp
then obtain tr1 s1' where ttr1-eq: ttr1 = tr1 ## s1'
using rev-exhaust by blast
hence tr1s1': never isIntO tr1 isIntO s1' using 13 by auto
have lfinite ltr1  $\implies$  s1'  $\neq$  llast ltr1

```

by (metis Opt.final-not-isInt Opt.lcompletedFrom-def llast-last-llist-of tr1s1'(2)
 vltr1(2))
 hence ne: lltr1' ≠ []
 using ltr1 unfolding ttr1-eq by auto
 then obtain s1'' ltr1' where lltr1': lltr1' = s1'' \$ ltr1'
 by (meson llist.exhaust)
 have [simp]: filter isIntO tr1 = []
 by (metis never-Nil-filter tr1s1'(1))
 have cltr1': Opt.lcompletedFrom s1 lltr1'
 by (metis Opt.lcompletedFrom-def lfinite-lappend lfinite-llist-of llast-lappend-LCons
 llast-last-llist-of lltr1' ltr1 ne vltr1(2))

 have inter4: ¬ lnever isIntO ltr2 using A inter3
 by (metis Opt.lA.eq-LNil-iff Opt.lO Opt.lO.eq-LNil-iff lfiltermap-LNil-never
 lfiltermap-lmap-lfilter vltr1(2) vltr2(2))
 have 04: ∃ s ∈ lset ltr2. isIntO s using inter4 unfolding llist.pred-set by auto
 define ttr2 where ttr2: ttr2 ≡ ltakeUntil isIntO ltr2
 define lltr2' where lltr2': lltr2' ≡ ldropUntil isIntO ltr2
 have ltr2: ltr2 = lappend (llist-of ttr2) lltr2'
 unfolding ttr2 lltr2' lappend-ltakeUntil-ldropUntil[OF 04] ..
 have 14: ttr2 ≠ [] ∧ never isIntO (butlast ttr2) ∧ isIntO (last ttr2)
 unfolding ttr2
 using ltakeUntil-last[OF 04] ltakeUntil-not-Nil[OF 04] ltakeUntil-never-butlast[OF
 04] by simp
 then obtain tr2 s2' where ttr2-eq: ttr2 = tr2 ## s2'
 using rev-exhaust by blast
 hence tr2s2': never isIntO tr2 isIntO s2' using 14 by auto
 have lfinite ltr2 ⇒ s2' ≠ llast ltr2
 by (metis Opt.final-not-isInt Opt.lcompletedFrom-def llast-last-llist-of tr2s2'(2)
 vltr2(2))
 hence ne: lltr2' ≠ []
 using ltr2 unfolding ttr2-eq by auto
 then obtain s2'' ltr2' where lltr2': lltr2' = s2'' \$ ltr2'
 by (meson llist.exhaust)
 have [simp]: filter isIntO tr2 = []
 by (metis never-Nil-filter tr2s2'(1))
 have cltr2': Opt.lcompletedFrom s2 lltr2'
 by (metis Opt.lcompletedFrom-def lfinite-lappend lfinite-llist-of llast-lappend-LCons
 llast-last-llist-of lltr2' ltr2 ne vltr2(2))

 have AA: Opt.lA lltr1' = Opt.lA lltr2'
 unfolding Opt.lA[OF cltr1'] Opt.lA[OF cltr2']
 using A[unfolded Opt.lA[OF vltr1(2)] Opt.lA[OF vltr2(2)]] tr1s1' tr2s2'
 unfolding ltr1 ltr2 ttr1-eq ttr2-eq
 unfolding lfilter-lappend-llist-of by simp

 show ?thesis apply(rule exI[of - tr1]) apply(rule exI[of - s1'])

apply(rule *exI*[of - *s1''*]) **apply**(rule *exI*[of - *ltr1'*])
apply(rule *exI*[of - *tr2*]) **apply**(rule *exI*[of - *s2'*])
apply(rule *exI*[of - *s2''*]) **apply**(rule *exI*[of - *ltr2'*])
unfolding $\varphi\varphi$ -def **apply**(intro conjI)
subgoal unfolding *ltr1 ttr1-eq lltr1'* ..
subgoal unfolding *ltr2 ttr2-eq lltr2'* ..
subgoal using *vltr1(1) unfolding ltr1 ttr1-eq lltr1'*
by (simp add: *Opt.lvalidFromS-lappend-finite lappend-llist-of-LCons*)
subgoal using *vltr2(1) unfolding ltr2 ttr2-eq lltr2'*
by (simp add: *Opt.lvalidFromS-lappend-finite lappend-llist-of-LCons*)
subgoal using *tr1s1'* **by** simp
subgoal using *tr2s2'* **by** simp
subgoal using *tr1s1'* **by** simp
subgoal using *tr2s2'* **by** simp
subgoal using $A[\text{unfolded } \text{Opt.lA}[OF \text{vltr1}(2)] \text{Opt.lA}[OF \text{vltr2}(2)]]$
tr1s1' tr2s2'
unfolding *ltr1 ttr1-eq ltr2 ttr2-eq lltr1' lltr2'*
unfolding *lfilter-lappend-llist-of* **by** simp
subgoal using *vltr1(1) unfolding ltr1 ttr1-eq lltr1'*
using *Opt.lvalidFromS-lappend-LCons* **by** blast
subgoal using *vltr1(2) unfolding ltr1 ttr1-eq lltr1'*
by (metis *Opt.lcompletedFrom-def lfinite-lappend lfinite-llist-of*
llast-lappend-LCons llast-last-llist-of llist.distinct(1))
subgoal using *vltr2(1) unfolding ltr2 ttr2-eq lltr2'*
using *Opt.lvalidFromS-lappend-LCons* **by** blast
subgoal using *vltr2(2) unfolding ltr2 ttr2-eq lltr2'*
by (metis *Opt.lcompletedFrom-def lfinite-lappend lfinite-llist-of*
llast-lappend-LCons llast-last-llist-of llist.distinct(1))
subgoal using AA **unfolding** *lltr1' lltr2'* ..

qed

definition $\chi\chi$ *s1 ltr1 tr1 s1' s1'' ltr1'* \equiv
ltr1 = lappend (llist-of (tr1 ## s1')) (s1'' \$ ltr1') \wedge
Opt.lvalidFromS s1 ((tr1 ## s1') ## s1'') \wedge
never isIntO tr1 \wedge \neg isIntO s1' \wedge \neg isIntO s1'' \wedge
never isSecO tr1 \wedge isSecO s1' \wedge
Opt.lvalidFromS s1'' (s1'' \$ ltr1') \wedge Opt.lcompletedFrom s1'' (s1'' \$ ltr1')

lemma *isSecO- $\chi\chi$* :

assumes *vltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1*

and *inter: lnever isIntO ltr1 and isec: \neg lnever isSecO ltr1*

shows $\exists tr1 s1' s1'' ltr1'. \chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$

proof –

have *0: \exists s \in lset ltr1. isSecO s* **using** *isec* **unfolding** *llist.pred-set* **by** auto

define *ttr1* **where** *ttr1: ttr1 \equiv ltakeUntil isSecO ltr1*

define *lltr1'* **where** *lltr1': lltr1' \equiv ldropUntil isSecO ltr1*

have *ltr1: ltr1 = lappend (llist-of ttr1) lltr1'*

unfolding *ttr1 lltr1' lappend-ltakeUntil-ldropUntil[OF 0]* ..

have *1: ttr1 \neq [] \wedge never isSecO (butlast ttr1) \wedge isSecO (last ttr1)*

```

unfolding ttr1
using ltakeUntil-last[OF 0] ltakeUntil-not-Nil[OF 0] ltakeUntil-never-butlast[OF 0]
by simp
then obtain tr1 s1' where ttr1-eq: ttr1 = tr1 ## s1'
using rev-exhaust by blast
hence tr1s1': never isSecO tr1 isSecO s1' using 1 by auto
have ?2: never isIntO tr1 ∧ ¬ isIntO s1' ∧ lnever isIntO lltr1'
using inter unfolding ltr1 ttr1-eq
unfolding llist-all-lappend-llist-of list-all-append by simp
have lfinite ltr1 ⇒ s1' ≠ llast ltr1
by (metis Opt.final-not-isSec Opt.lcompletedFrom-def llast-last-llist-of tr1s1'?(2)
vltr1?(2))
hence ne: lltr1' ≠ []
using ltr1 unfolding ttr1-eq by auto
then obtain s1'' ltr1' where lltr1': lltr1' = s1'' $ ltr1'
by (meson llist.exhaust)
show ?thesis apply(rule exI[of - tr1]) apply(rule exI[of - s1'])
apply(rule exI[of - s1'']) apply(rule exI[of - ltr1'])
unfolding χχ-def apply(intro conjI)
  subgoal unfolding ltr1 ttr1-eq lltr1' ..
  subgoal using vltr1?(1) unfolding ltr1 ttr1-eq lltr1'
  by (simp add: Opt.lvalidFromS-lappend-finite lappend-llist-of-LCons)
  subgoal using ?2 by simp
  subgoal using ?2 by simp
  subgoal using ?2 unfolding lltr1' by simp
  subgoal using tr1s1' by simp
  subgoal using tr1s1' by simp
  subgoal using vltr1?(1) unfolding ltr1 ttr1-eq lltr1'
  using Opt.lvalidFromS-lappend-LCons by blast
  subgoal using vltr1?(2) unfolding ltr1 ttr1-eq lltr1'
  by (metis Opt.lcompletedFrom-def ne lfinite-lappend
lfinite-llist-of llast-lappend-LCons llast-last-llist-of lltr1') .
qed

```

```

type-synonym ('stA, 'stO) tuple34 =
  enat × enat ×
  'stA × 'stA llist ×
  'stA × 'stA llist ×
  status ×
  'stO × 'stO × status

```

```

type-synonym ('stA, 'stO) tuple12 =
  'stO list × 'stO × 'stO list × 'stO × status × status

```

context

fixes $\Delta :: \text{enat} \Rightarrow \text{enat} \Rightarrow \text{enat} \Rightarrow \text{'stateO} \Rightarrow \text{'stateO} \Rightarrow \text{status} \Rightarrow \text{'stateV} \Rightarrow \text{'stateV} \Rightarrow \text{status} \Rightarrow \text{bool}$

begin

fun $\text{isn} :: \text{turn} \times (\text{'stateO}, \text{'stateV}) \text{tuple34} \Rightarrow \text{bool}$

where

$\text{isn} (\text{trn}, w1, w2, s1, \text{ltr1}, s2, \text{ltr2}, \text{statA}, sv1, sv2, \text{statO}) \longleftrightarrow \text{ltr1} = [] \wedge \text{ltr2} = []$

fun $\text{h-t} ::$

$\text{turn} \times (\text{'stateO}, \text{'stateV}) \text{tuple34} \Rightarrow$

$(\text{'stateO}, \text{'stateV}) \text{tuple12} \times$

$\text{turn} \times (\text{'stateO}, \text{'stateV}) \text{tuple34}$

where

$\text{h-t} (\text{trn}, w1, w2, s1, \text{ltr1}, s2, \text{ltr2}, \text{statA}, sv1, sv2, \text{statO}) =$

$(\text{if } \text{trn} = L$

$\text{then if } \text{lnever } \text{isSecO } \text{ltr1}$

$\text{then let } (s1', \text{ltr1}') = (\text{lhd } (\text{ttl } \text{ltr1}), \text{ttl } \text{ltr1})$

$\text{in let } (w1', w2', \text{trv1}, sv1', \text{trv2}, sv2', \text{statOO}) =$

$(\text{SOME } k. \text{ case } k \text{ of } (w1', w2', \text{trv1}, sv1', \text{trv2}, sv2', \text{statOO}) \Rightarrow$

$\omega3 \Delta w1 w2 w1' w2' s1 s1' s2 \text{statA } sv1 \text{trv1 } sv1' sv2 \text{trv2 } sv2' \text{statOO})$

$\text{in } ((\text{trv1}, sv1', \text{trv2}, sv2', \text{statA}, \text{statOO}),$

$(\text{if } \text{trv1} = [] \text{ then } L \text{ else } R,$

$w1', w2', s1', \text{ltr1}', s2, \text{ltr2}, \text{statA}, sv1', sv2', \text{statOO}))$

else

$\text{let } (\text{tr1}, s1', s1'', \text{ltr1}') =$

$(\text{SOME } k. \text{ case } k \text{ of } (\text{tr1}, s1', s1'', \text{ltr1}') \Rightarrow$

$\chi\chi s1 \text{ltr1 } \text{tr1 } s1' s1'' \text{ltr1}')$

$\text{in let } (w1', w2', \text{trv1}, sv1'', \text{trv2}, sv2'', \text{statOO}) =$

$(\text{SOME } k'. \text{ case } k' \text{ of } (w1', w2', \text{trv1}, sv1'', \text{trv2}, sv2'', \text{statOO}) \Rightarrow$

$\chi3' \Delta w1 w2 w1' w2' s1 \text{tr1 } s1' s1'' s2 \text{statA } sv1 \text{trv1 } sv1'' sv2 \text{trv2}$

$sv2'' \text{statOO})$

$\text{in } ((\text{trv1}, sv1'', \text{trv2}, sv2'', \text{statA}, \text{statOO}),$

$(R, w1', w2', s1'', s1'' \$ \text{ltr1}', s2, \text{ltr2}, \text{statA}, sv1'', sv2'', \text{statOO}))$

$\text{else if } \text{lnever } \text{isSecO } \text{ltr2}$

$\text{then let } (s2', \text{ltr2}') = (\text{lhd } (\text{ttl } \text{ltr2}), \text{ttl } \text{ltr2})$

$\text{in let } (w1', w2', \text{trv1}, sv1', \text{trv2}, sv2', \text{statOO}) =$

$(\text{SOME } k. \text{ case } k \text{ of } (w1', w2', \text{trv1}, sv1', \text{trv2}, sv2', \text{statOO}) \Rightarrow$

$\omega4 \Delta w1 w2 w1' w2' s1 s2 s2' \text{statA } sv1 \text{trv1 } sv1' sv2 \text{trv2 } sv2' \text{statOO})$

$\text{in } ((\text{trv1}, sv1', \text{trv2}, sv2', \text{statA}, \text{statOO}),$

$(\text{if } \text{trv2} = [] \text{ then } R \text{ else } L,$

$w1', w2', s1, \text{ltr1}, s2', \text{ltr2}', \text{statA}, sv1', sv2', \text{statOO}))$

else

$\text{let } (\text{tr2}, s2', s2'', \text{ltr2}') =$

$(\text{SOME } k. \text{ case } k \text{ of } (\text{tr2}, s2', s2'', \text{ltr2}') \Rightarrow$

$\chi\chi s2 \text{ltr2 } \text{tr2 } s2' s2'' \text{ltr2}')$

$\text{in let } (w1', w2', \text{trv1}, sv1'', \text{trv2}, sv2'', \text{statOO}) =$

```

      (SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statOO) =>
         $\chi_4'$   $\Delta$  w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2
        sv2'' statOO)
      in ((trv1,sv1'',trv2,sv2'',statA,statOO),
        (L,w1',w2',s1, ltr1, s2'',s2'' $ ltr2',statA,sv1'',sv2'',statOO))
    )

```

declare *h-t.simps*[simp del]

definition *h* \equiv *fst o h-t*

definition *t* \equiv *snd o h-t*

fun *econd* **where** *econd* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) =
 (*llength ltr1* \leq *Suc 0* \vee *llength ltr2* \leq *Suc 0*)

fun *e* **where** *e* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) = [[([*sv1*],*sv1*],[*sv2*],*sv2,statA,statO*)]]

definition *f* :: *turn* \times ('*stateO*,'*stateV*)*tuple34* \Rightarrow ('*stateO*,'*stateV*)*tuple12* *llist*
where *f* \equiv *ccorec-llist isn h econd e t*

lemma *f-LNil*:

ltr1 = [] \Rightarrow *ltr2* = [] \Rightarrow *f* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) =
 []

unfolding *f-def* **apply**(*subst llist-ccorec(1)*) **by** *auto*

lemma *f-length-1*:

assumes *ltr1* \neq [] \vee *ltr2* \neq [] *llength ltr1* \leq *Suc 0* \vee *llength ltr2* \leq *Suc 0*

shows *f* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) = [[([*sv1*],*sv1*],[*sv2*],*sv2,statA,statO*)]]

using *assms* **unfolding** *f-def* **apply**(*subst llist-ccorec(2)*)

subgoal **unfolding** *e.simps lnull-def* **by** *auto*

subgoal **by** *auto*

subgoal **unfolding** *econd.simps* **by** *simp* .

lemma *f-length-ge1*:

assumes *llength ltr1* $>$ *Suc 0* *llength ltr2* $>$ *Suc 0*

shows *f* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*) =

LCons (*h* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*)) (*f* (*t* (*trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO*))

proof –

show *?thesis* **using** *assms* **unfolding** *f-def* **apply**(*subst llist-ccorec(2)*)

subgoal **unfolding** *e.simps lnull-def* **by** *auto*

subgoal **by** *auto*

subgoal **unfolding** *econd.simps* **by** *auto* .

qed

definition $lltrv1 :: turn \times ('stateO, 'stateV)tuple3_4 \Rightarrow 'stateV\ llist$ **where**
 $lltrv1\ trn\text{-}tp = lconcat\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO).\ llist\text{-}of\ trv1)\ (f\ trn\text{-}tp))$

definition $then1 :: turn \times ('stateO, 'stateV)tuple3_4 \Rightarrow nat$ **where**
 $then1\ trn\text{-}tp = firstNC\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO).\ trv1)\ (f\ trn\text{-}tp))$

definition $lltrv2 :: turn \times ('stateO, 'stateV)tuple3_4 \Rightarrow 'stateV\ llist$ **where**
 $lltrv2\ trn\text{-}tp = lconcat\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO).\ llist\text{-}of\ trv2)\ (f\ trn\text{-}tp))$

definition $then2 :: turn \times ('stateO, 'stateV)tuple3_4 \Rightarrow nat$ **where**
 $then2\ trn\text{-}tp = firstNC\ (lmap\ (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO).\ trv2)\ (f\ trn\text{-}tp))$

lemma $lltrv1\text{-}ne\text{-}imp$:

assumes $lltrv1\ trn\text{-}tp \neq []$

shows $\exists trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO.\ (trv1, sv1'', trv2, sv2'', statAA, statOO) \in lset\ (f\ trn\text{-}tp) \wedge trv1 \neq []$

using $assms\ unfolding\ lltrv1\text{-}def\ unfolding\ lconcat\text{-}eq\ LNil\text{-}iff$ **by** $force$

lemma $lltrv2\text{-}ne\text{-}imp$:

assumes $lltrv2\ trn\text{-}tp \neq []$

shows $\exists trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO.\ (trv1, sv1'', trv2, sv2'', statAA, statOO) \in lset\ (f\ trn\text{-}tp) \wedge trv2 \neq []$

using $assms\ unfolding\ lltrv2\text{-}def\ unfolding\ lconcat\text{-}eq\ LNil\text{-}iff$ **by** $force$

lemma $lltrv1\text{-}LNil[simp]$:

$ltr1 = [] \implies ltr2 = [] \implies lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = []$

unfolding $lltrv1\text{-}def\ f\text{-}LNil$ **by** $simp$

lemma $lltrv2\text{-}LNil[simp]$:

$ltr1 = [] \implies ltr2 = [] \implies lltrv2\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = []$

unfolding $lltrv2\text{-}def\ f\text{-}LNil$ **by** $simp$

lemma $lltrv1\text{-}lnever[simp]$:

assumes $ltr1 \neq [] \vee ltr2 \neq [] \wedge llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$

shows $lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = [[sv1]]$

unfolding $lltrv1\text{-}def$ **using** $f\text{-}length\text{-}1[OF\ assms]$ **by** $auto$

lemma *lltrv2-lnever[simp]*:
assumes $ltr1 \neq [] \vee ltr2 \neq []$ $llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$
shows $lltrv2\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = [[sv2]]$
unfolding *lltrv2-def* **using** *f-length-1[OF assms]* **by** *auto*

lemma *h-t-lnever-L*:
assumes *unw*: *unwindCond* Δ
and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1*
and *l'*: *lnever isIntO* *ltr1* \neg *isIntO* *s2*
and *len*: $llength\ ltr1 > Suc\ 0$ $llength\ ltr2 > Suc\ 0$
and *l*: $trn = L$ *lnever isSecO* *ltr1*
shows $\exists w1'\ w2'\ s1'\ ltr1'\ trv1\ sv1'\ trv2\ sv2'\ statOO.$
 $ltr1 = s1\ \$\ ltr1' \wedge validTransO\ (s1, s1') \wedge$
 $Opt.lvalidFromS\ s1'\ ltr1' \wedge Opt.lcompletedFrom\ s1'\ ltr1' \wedge lnever\ isIntO\ ltr1' \wedge$

$\omega 3\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s1'\ s2\ statA\ sv1\ trv1\ sv1'\ sv2\ trv2\ sv2'\ statOO \wedge$
 $h-t\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1', trv2, sv2', statA, statOO),$
 $(if\ trv1 = []\ then\ L\ else\ R,$
 $w1', w2', s1', ltr1', s2, ltr2, statA, sv1', sv2', statOO))$

proof–
have *s1*: $\neg\ isIntO\ s1$ **using** *l'* *ltr1*
by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *lhd-LCons* *llist.exhaust* *llist.pred-inject*(2))

obtain *ltr1'* **where** *ltr13*: $ltr1 = s1\ \$\ ltr1'$
by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *llist.exhaust-sel* *ltr1*(1) *ltr1*(2))
hence *ltr1'*: $ltr1' = ltl\ ltr1$ **by** *auto*
have *ltr1'ne*: $ltr1' \neq []$ **using** *len*(1) **unfolding** *ltr13*
by (*metis* *One-nat-def* $llength-LCons$ $llength-LNil$ *one-eSuc* *one-enat-def* *order-less-irrefl*)
define *s1'* **where** *s1'*: $s1' = lhd\ (ltl\ ltr1)$
have *v3*: $validTransO\ (s1, s1')$ **and** *vv3*: *Opt.lvalidFromS* *s1'* *ltr1'* *Opt.lcompletedFrom* *s1'* *ltr1'*
using *ltr1* *ltr1'ne* **unfolding** *ltr13* *s1'*
by (*metis* *Opt.lcompletedFrom-LCons* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-Cons-iff* *ltr1'* *ltr13*)

have *is1'*: $\neg\ isIntO\ s1'$ **and** *lnever isIntO* *ltr1'*
using *l'*(1) **unfolding** *ltr13*
by (*metis* *llist.exhaust-sel* *llist.pred-inject*(2) *ltr1'* *ltr1'ne* *s1'*)
+


```

have iss1:  $\neg$  isSecO s1
  using l(2) ltr13 by auto

obtain w1' w2' trv1 sv1' trv2 sv2' statOO
where  $\omega3$ :  $\omega3 \Delta w1 w2 w1' w2' s1$  (lhd (ltl ltr1)) s2 statA sv1 trv1 sv1' sv2
trv2 sv2' statOO
using unwindCond-ex- $\omega3$ [OF unW  $\Delta$  r v3 s1 is1' iss1 l'(2)] s1' by auto

define tp' where
tp' = (SOME k'. case k' of (w1',w2',trv1,sv1',trv2,sv2',statOO)  $\Rightarrow$ 
   $\omega3 \Delta w1 w2 w1' w2' s1$  (lhd (ltl ltr1)) s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO)

have 1: case tp' of (w1',w2',trv1,sv1',trv2,sv2',statOO)  $\Rightarrow$ 
   $\omega3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2' statOO$ 
using  $\omega3$  unfolding tp'-def s1' apply– apply(rule someI-ex)
apply(rule exI[of - (w1',w2',trv1,sv1',trv2,sv2',statOO)]) by auto

obtain w1' w2' trv1 sv1' trv2 sv2' statOO where
tp': tp' = (w1',w2',trv1,sv1',trv2,sv2',statOO) by(cases tp', auto)

have  $\omega3$ :  $\omega3 \Delta w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2' statOO$ 

using 1 unfolding tp' by auto

show ?thesis
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1'])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using len l unfolding h-t.simps apply simp
  unfolding tp'-def[symmetric] tp' s1' ltr1' by simp .
qed

lemma lltrv1-lltrv2-lnever-L:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1

```

and l' : $\text{lnever isIntO } ltr1 \neg \text{isIntO } s2$
and len : $\text{llength } ltr1 > \text{Suc } 0 \text{ llength } ltr2 > \text{Suc } 0$
and l : $trn = L \text{ lnever isSecO } ltr1$
shows $\exists w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' \text{ statOO}$.
 $ltr1 = s1 \ \$ \ ltr1' \wedge \text{validTransO } (s1, s1') \wedge$
 $\text{Opt.lvalidFromS } s1' \ ltr1' \wedge \text{Opt.lcompletedFrom } s1' \ ltr1' \wedge \text{lnever isIntO } ltr1' \wedge$
 $\omega3 \ \Delta \ w1 \ w2 \ w1' \ w2' \ s1 \ s1' \ s2 \ \text{statA } sv1 \ trv1 \ sv1' \ sv2 \ trv2 \ sv2' \ \text{statOO} \wedge$
 $\text{lltrv1 } (trn, w1, w2, s1, ltr1, s2, ltr2, \text{statA}, sv1, sv2, \text{statO}) =$
 $\text{lappend } (\text{llist-of } trv1) \ (\text{lltrv1 } (\text{if } trv1 = [] \text{ then } L \text{ else } R,$
 $\quad w1', w2', s1', ltr1', s2, ltr2, \text{statA}, sv1', sv2', \text{statOO})) \ \wedge$
 $\text{lltrv2 } (trn, w1, w2, s1, ltr1, s2, ltr2, \text{statA}, sv1, sv2, \text{statO}) =$
 $\text{lappend } (\text{llist-of } trv2) \ (\text{lltrv2 } (\text{if } trv1 = [] \text{ then } L \text{ else } R,$
 $\quad w1', w2', s1', ltr1', s2, ltr2, \text{statA}, sv1', sv2', \text{statOO}))$

proof –

show *?thesis*

using $h\text{-t-lnever-L}[OF \ \text{assms}]$ **apply**($\text{elim } exE$)

subgoal for $w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' \text{ statOO}$

apply($\text{rule } exI[\text{of } - \ w1']$) **apply**($\text{rule } exI[\text{of } - \ w2']$)

apply($\text{rule } exI[\text{of } - \ s1']$) **apply**($\text{rule } exI[\text{of } - \ ltr1']$)

apply($\text{rule } exI[\text{of } - \ trv1]$) **apply**($\text{rule } exI[\text{of } - \ sv1']$)

apply($\text{rule } exI[\text{of } - \ trv2]$) **apply**($\text{rule } exI[\text{of } - \ sv2']$)

apply($\text{rule } exI[\text{of } - \ \text{statOO}]$)

apply($\text{intro } conjI$)

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal unfolding lltrv1-def **apply**($\text{subst } f\text{-length-ge1}[OF \ len]$)

unfolding $h\text{-def } t\text{-def}$ **by** *auto*

subgoal unfolding lltrv2-def **apply**($\text{subst } f\text{-length-ge1}[OF \ len]$)

unfolding $h\text{-def } t\text{-def}$ **by** *auto* . .

qed

lemma $h\text{-t-not-lnever-L}$:

assumes unw : $\text{unwindCond } \Delta$

and Δ : $\Delta \ \infty \ w1 \ w2 \ s1 \ s2 \ \text{statA } sv1 \ sv2 \ \text{statO}$

and r : $\text{reachO } s1 \ \text{reachO } s2 \ \text{reachV } sv1 \ \text{reachV } sv2$

and $ltr1$: $\text{Opt.lvalidFromS } s1 \ ltr1 \ \text{Opt.lcompletedFrom } s1 \ ltr1$

and l' : $\text{lnever isIntO } ltr1 \neg \text{isIntO } s2$

and len : $\text{llength } ltr1 > \text{Suc } 0 \text{ llength } ltr2 > \text{Suc } 0$

and l : $trn = L \ \neg \text{lnever isSecO } ltr1$

shows $\exists w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' \text{ statOO}$.

$\chi\chi \ s1 \ ltr1 \ tr1 \ s1' \ s1'' \ ltr1' \ \wedge$

$\chi\chi^3 \ \Delta \ w1 \ w2 \ w1' \ w2' \ s1 \ tr1 \ s1' \ s1'' \ s2 \ \text{statA } sv1 \ trv1 \ sv1'' \ sv2 \ trv2 \ sv2'' \ \text{statOO}$

\wedge
 $h\text{-}t (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1'', trv2, sv2'', statA, statOO),$
 $(R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO))$

proof –

have $s1: \neg isIntO s1$ **using** $l' ltr1$
by $(metis Opt.lvalidFromS-def l(2) lhd-LCons llist.exhaust llist.pred-inject(1)$
 $l\text{list}.pred\text{-}inject(2))$

obtain $tr1 s1' s1'' ltr1'$
where $\chi\chi: \chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$
using $isSecO\text{-}\chi\chi[OF ltr1 l'(1) l(2)]$ **by** $auto$

define tp **where**
 $tp = (SOME k. case k of (tr1, s1', s1'', ltr1') \Rightarrow$
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1')$

have $0: case tp of (tr1, s1', s1'', ltr1') \Rightarrow$
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$
using $\chi\chi$ **unfolding** $tp\text{-}def$ **apply– apply** $(rule someI\text{-}ex)$
apply $(rule exI[of - (tr1, s1', s1'', ltr1')])$ **by** $auto$

obtain $tr1 s1' s1'' ltr1'$ **where**
 $tp: tp = (tr1, s1', s1'', ltr1')$ **by** $(cases tp, auto)$

have $\chi\chi: \chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$
using 0 **unfolding** tp **by** $auto$

obtain $w1' w2' trv1 sv1'' trv2 sv2'' statOO$
where $\chi^3: \chi^3 \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2$
 $sv2'' statOO$
using $unwindCond\text{-}ex\text{-}\chi^3[OF unw \Delta r, of tr1 s1' s1'']$
using $\chi\chi l' s1$ **unfolding** $\chi\chi\text{-}def$ **by** $auto$

define tp' **where**
 $tp' = (SOME k'. case k' of (w1', w2', trv1, sv1'', trv2, sv2'', statOO) \Rightarrow$
 $\chi^3 \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2''$
 $statOO)$

have $1: case tp' of (w1', w2', trv1, sv1'', trv2, sv2'', statOO) \Rightarrow$
 $\chi^3 \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2''$
 $statOO$
using χ^3 **unfolding** $tp'\text{-}def$ **apply– apply** $(rule someI\text{-}ex)$
apply $(rule exI[of - (w1', w2', trv1, sv1'', trv2, sv2'', statOO)])$ **by** $auto$

obtain $w1' w2' trv1 sv1'' trv2 sv2'' statOO$ **where**
 $tp': tp' = (w1', w2', trv1, sv1'', trv2, sv2'', statOO)$ **by** $(cases tp', auto)$

```

have  $\chi^3$ :  $\chi^3 \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2$ 
 $sv2'' statOO$ 
using 1 unfolding  $tp'$  by auto

show ?thesis
apply(rule  $exI[of - w1]$ ) apply(rule  $exI[of - w2]$ )
apply(rule  $exI[of - tr1]$ ) apply(rule  $exI[of - s1]$ ) apply(rule  $exI[of - s1']$ )
apply(rule  $exI[of - ltr1]$ )
apply(rule  $exI[of - trv1]$ ) apply(rule  $exI[of - sv1']$ ) apply(rule  $exI[of - trv2]$ )
apply(rule  $exI[of - sv2']$ )
apply(rule  $exI[of - statOO]$ )
apply(intro conjI)
subgoal using  $\chi\chi$  .
subgoal using  $\chi^3$  .
subgoal using l unfolding h-t.simps
unfolding  $tp-def[symmetric]$   $tp$  apply simp
unfolding  $tp'-def[symmetric]$   $tp'$  by simp .
qed

```

```

lemma lltrv1-lltrv2-not-lnever-L:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and  $r$ : reachO  $s1$  reachO  $s2$  reachV  $sv1$  reachV  $sv2$ 
and  $ltr1$ : Opt.lvalidFromS  $s1$   $ltr1$  Opt.lcompletedFrom  $s1$   $ltr1$ 
and  $l'$ : lnever isIntO  $ltr1$   $\neg$  isIntO  $s2$ 
and  $len$ : llength  $ltr1$   $>$  Suc 0 llength  $ltr2$   $>$  Suc 0
and  $l$ :  $trn = L$   $\neg$  lnever isSecO  $ltr1$ 
shows  $\exists w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO.$ 
 $\chi\chi s1 ltr1 tr1 s1' s1'' ltr1' \wedge$ 
 $\chi^3 \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
 $\wedge$ 
lltrv1 ( $trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO$ ) =
lappend (llist-of  $trv1$ ) (lltrv1 ( $R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO$ ))
 $\wedge$ 
lltrv2 ( $trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO$ ) =
lappend (llist-of  $trv2$ ) (lltrv2 ( $R, w1', w2', s1'', s1'' \$ ltr1', s2, ltr2, statA, sv1'', sv2'', statOO$ ))

```

```

proof –
show ?thesis
using h-t-not-lnever-L[OF assms] apply(elim exE)
subgoal for  $w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO$ 
apply(rule  $exI[of - w1]$ ) apply(rule  $exI[of - w2]$ )
apply(rule  $exI[of - tr1]$ ) apply(rule  $exI[of - s1]$ ) apply(rule  $exI[of - s1']$ )
apply(rule  $exI[of - ltr1]$ )
apply(rule  $exI[of - trv1]$ ) apply(rule  $exI[of - sv1']$ ) apply(rule  $exI[of - trv2]$ )
apply(rule  $exI[of - sv2']$ )
apply(rule  $exI[of - statOO]$ )
apply(intro conjI)

```

subgoal by simp
subgoal by simp
subgoal unfolding ltrv1-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by simp
subgoal unfolding ltrv2-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by simp . .
qed

lemma h-t-lnever-R:

assumes unw: unwindCond Δ
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and l': $\neg isIntO s1 lnever isIntO ltr2$
and len: llength ltr1 > Suc 0 llength ltr2 > Suc 0
and l: trn = R lnever isSecO ltr2
shows $\exists w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO.$
 $ltr2 = s2 \$ ltr2' \wedge validTransO (s2, s2') \wedge$
 $Opt.lvalidFromS s2' ltr2' \wedge Opt.lcompletedFrom s2' ltr2' \wedge lnever isIntO ltr2' \wedge$

$\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2' statOO \wedge$
 $h-t (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1', trv2, sv2', statA, statOO),$
 $(if trv2 = [] then R else L,$
 $w1', w2', s1, ltr1, s2', ltr2', statA, sv1', sv2', statOO))$

proof –

have s2: $\neg isIntO s2$ using l' ltr2
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil lhd-LCons
 $l\list.exhaust l\list.pred-inject(2))$

obtain ltr2' where ltr24: $ltr2 = s2 \$ ltr2'$
by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil l\list.exhaust-sel
 $ltr2(1) ltr2(2))$
hence ltr2': $ltr2' = ltl ltr2$ by auto
have ltr2'ne: $ltr2' \neq []$ using len(2) unfolding ltr24
by (metis One-nat-def llength-LCons llength-LNil one-eSuc one-enat-def or-
 $der-less-irrefl)$
define s2' where s2': $s2' = lhd (ltl ltr2)$
have v4: $validTransO (s2, s2')$ and vv4: $Opt.lvalidFromS s2' ltr2' Opt.lcompletedFrom$
 $s2' ltr2'$
using ltr2 ltr2'ne unfolding ltr24 s2'
by (metis Opt.lcompletedFrom-LCons Opt.lcompletedFrom-def Opt.lvalidFromS-Cons-iff
 $ltr2' ltr24)+$

have is2': $\neg isIntO s2'$ and lnever isIntO ltr2'
using l'(2) unfolding ltr24
by (metis l\list.exhaust-sel l\list.pred-inject(2) ltr2' ltr2'ne s2')+

```

have iss2:  $\neg$  isSecO s2
  using l(2) ltr24 by auto

obtain w1' w2' trv1 sv1' trv2 sv2' statOO
  where  $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2$  (lhd (ltl ltr2)) statA sv1 trv1 sv1' sv2
trv2 sv2' statOO
  using unwindCond-ex- $\omega_4$ [OF unw  $\Delta$  r l'(1) v4 s2 is2' iss2] s2' by auto

define tp' where
  tp' = (SOME k'. case k' of (w1',w2',trv1,sv1',trv2,sv2',statOO)  $\Rightarrow$ 
     $\omega_4 \Delta w1 w2 w1' w2' s1 s2$  (lhd (ltl ltr2)) statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO)

  have 1: case tp' of (w1',w2',trv1,sv1',trv2,sv2',statOO)  $\Rightarrow$ 
     $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2'$  statA sv1 trv1 sv1' sv2 trv2 sv2' statOO
  using  $\omega_4$  unfolding tp'-def s2' apply– apply(rule someI-ex)
  apply(rule exI[of - (w1',w2',trv1,sv1',trv2,sv2',statOO)]) by auto

obtain w1' w2' trv1 sv1' trv2 sv2' statOO where
  tp': tp' = (w1',w2',trv1,sv1',trv2,sv2',statOO) by(cases tp', auto)

have  $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2'$  statA sv1 trv1 sv1' sv2 trv2 sv2' statOO

using 1 unfolding tp' by auto

show ?thesis
apply(rule exI[of - w1  $\uparrow$ ]) apply(rule exI[of - w2  $\uparrow$ ])
apply(rule exI[of - s2  $\uparrow$ ]) apply(rule exI[of - ltr2  $\uparrow$ ])
apply(rule exI[of - trv1]) apply(rule exI[of - sv1  $\uparrow$ ])
apply(rule exI[of - trv2]) apply(rule exI[of - sv2  $\uparrow$ ])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using len l unfolding h-t.simps apply simp
  unfolding tp'-def[symmetric] tp' s2' ltr2' by simp .
qed

lemma lltrv1-lltrv2-lnever-R:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2$  statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2

```

and $ltr2$: $Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$
and l' : $\neg\ isIntO\ s1\ lnever\ isIntO\ ltr2$
and len : $llength\ ltr1\ >\ Suc\ 0\ llength\ ltr2\ >\ Suc\ 0$
and l : $trn = R\ lnever\ isSecO\ ltr2$
shows $\exists\ w1'\ w2'\ s2'\ ltr2'\ trv1\ sv1'\ trv2\ sv2'\ statOO.$
 $ltr2 = s2\ \$\ ltr2' \wedge\ validTransO\ (s2, s2') \wedge$
 $Opt.lvalidFromS\ s2'\ ltr2' \wedge\ Opt.lcompletedFrom\ s2'\ ltr2' \wedge\ lnever\ isIntO\ ltr2' \wedge$

$\omega_4\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s2\ s2'\ statA\ sv1\ trv1\ sv1'\ sv2\ trv2\ sv2'\ statOO \wedge$
 $lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend\ (l\ list\ of\ trv1)\ (lltrv1\ (if\ trv2 = []\ then\ R\ else\ L,$
 $w1', w2', s1, ltr1, s2', ltr2', statA, sv1', sv2', statOO)) \wedge$
 $lltrv2\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend\ (l\ list\ of\ trv2)\ (lltrv2\ (if\ trv2 = []\ then\ R\ else\ L,$
 $w1', w2', s1, ltr1, s2', ltr2', statA, sv1', sv2', statOO))$

proof –

show *?thesis*

using $h\text{-}t\text{-}lnever\text{-}R[OF\ assms]$ **apply**($elim\ exE$)

subgoal for $w1'\ w2'\ s2'\ ltr2'\ trv1\ sv1'\ trv2\ sv2'\ statOO$

apply($rule\ exI[of\ -\ w1\ \uparrow]$) **apply**($rule\ exI[of\ -\ w2\ \uparrow]$)

apply($rule\ exI[of\ -\ s2\ \uparrow]$) **apply**($rule\ exI[of\ -\ ltr2\ \uparrow]$)

apply($rule\ exI[of\ -\ trv1]$) **apply**($rule\ exI[of\ -\ sv1\ \uparrow]$)

apply($rule\ exI[of\ -\ trv2]$) **apply**($rule\ exI[of\ -\ sv2\ \uparrow]$)

apply($rule\ exI[of\ -\ statOO]$)

apply($intro\ conjI$)

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal by *simp*

subgoal unfolding $lltrv1\text{-}def$ **apply**($subst\ f\text{-}length\text{-}ge1[OF\ len]$)

unfolding $h\text{-}def\ t\text{-}def$ **by** *auto*

subgoal unfolding $lltrv2\text{-}def$ **apply**($subst\ f\text{-}length\text{-}ge1[OF\ len]$)

unfolding $h\text{-}def\ t\text{-}def$ **by** *auto* . .

qed

lemma $h\text{-}t\text{-}not\text{-}lnever\text{-}R$:

assumes unw : $unwindCond\ \Delta$

and Δ : $\Delta\ \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and r : $reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$

and $ltr2$: $Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$

and l' : $\neg\ isIntO\ s1\ lnever\ isIntO\ ltr2$

and len : $llength\ ltr1\ >\ Suc\ 0\ llength\ ltr2\ >\ Suc\ 0$

and l : $trn = R\ \neg\ lnever\ isSecO\ ltr2$

shows $\exists\ w1'\ w2'\ tr2\ s2'\ s2''\ ltr2'\ trv1\ sv1''\ trv2\ sv2''\ statOO.$

$\chi\chi\ s2\ ltr2\ tr2\ s2'\ s2''\ ltr2' \wedge$

```

 $\chi_4' \Delta w_1 w_2 w_1' w_2' s_1 s_2 tr_2 s_2' s_2'' statA sv_1 trv_1 sv_1'' sv_2 trv_2 sv_2'' statOO$ 
 $\wedge$ 
 $h-t (trn, w_1, w_2, s_1, ltr_1, s_2, ltr_2, statA, sv_1, sv_2, statO) =$ 
 $((trv_1, sv_1'', trv_2, sv_2'', statA, statOO),$ 
 $(L, w_1', w_2', s_1, ltr_1, s_2'', s_2'' \$ ltr_2', statA, sv_1'', sv_2'', statOO))$ 
proof –
  have  $s_2: \neg isIntO s_2$  using  $l' ltr_2$ 
  by (metis Simple-Transition-System.lvalidFromS-def  $l(2)$  lhd-LCons llist.pred-inject(1)

    llist.pred-inject(2) neq-LNil-conv)

  obtain  $tr_2 s_2' s_2'' ltr_2'$ 
  where  $\chi\chi: \chi\chi s_2 ltr_2 tr_2 s_2' s_2'' ltr_2'$ 
  using isSecO- $\chi\chi$ [OF ltr_2 l'(2) l(2)] by auto

  define  $tp$  where
   $tp = (SOME k. case k of (tr_2, s_2', s_2'', ltr_2') \Rightarrow$ 
     $\chi\chi s_2 ltr_2 tr_2 s_2' s_2'' ltr_2')$ 

  have  $0: case tp of (tr_2, s_2', s_2'', ltr_2') \Rightarrow$ 
     $\chi\chi s_2 ltr_2 tr_2 s_2' s_2'' ltr_2'$ 
  using  $\chi\chi$  unfolding  $tp-def$  apply– apply(rule someI-ex)
  apply(rule exI[of - (tr_2, s_2', s_2'', ltr_2')]) by auto

  obtain  $tr_2 s_2' s_2'' ltr_2'$  where
   $tp: tp = (tr_2, s_2', s_2'', ltr_2')$  by(cases tp, auto)

  have  $\chi\chi: \chi\chi s_2 ltr_2 tr_2 s_2' s_2'' ltr_2'$ 
  using  $0$  unfolding  $tp$  by auto

  obtain  $w_1' w_2' trv_1 sv_1'' trv_2 sv_2'' statOO$ 
  where  $\chi_4': \chi_4' \Delta w_1 w_2 w_1' w_2' s_1 s_2 tr_2 s_2' s_2'' statA sv_1 trv_1 sv_1'' sv_2 trv_2$ 
 $sv_2'' statOO$ 
  using unwindCond-ex- $\chi_4'$ [OF unw  $\Delta r$ , of tr_2 s_2' s_2']
  using  $\chi\chi l' s_2$  unfolding  $\chi\chi-def$  by auto

  define  $tp'$  where
   $tp' = (SOME k'. case k' of (w_1', w_2', trv_1, sv_1'', trv_2, sv_2'', statOO) \Rightarrow$ 
     $\chi_4' \Delta w_1 w_2 w_1' w_2' s_1 s_2 tr_2 s_2' s_2'' statA sv_1 trv_1 sv_1'' sv_2 trv_2 sv_2''$ 
 $statOO)$ 

  have  $1: case tp' of (w_1', w_2', trv_1, sv_1'', trv_2, sv_2'', statOO) \Rightarrow$ 
     $\chi_4' \Delta w_1 w_2 w_1' w_2' s_1 s_2 tr_2 s_2' s_2'' statA sv_1 trv_1 sv_1'' sv_2 trv_2 sv_2''$ 
 $statOO$ 
  using  $\chi_4'$  unfolding  $tp'-def$  apply– apply(rule someI-ex)
  apply(rule exI[of - (w_1', w_2', trv_1, sv_1'', trv_2, sv_2'', statOO)]) by auto

```


obtain $w1' w2' trv1 sv1'' trv2 sv2'' statOO$ **where**
 tp' : $tp' = (w1', w2', trv1, sv1'', trv2, sv2'', statOO)$ **by**(cases tp' , auto)

have $\chi4'$: $\chi4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
using 1 **unfolding** tp' **by** auto

show ?thesis
apply(rule $exI[of - w1']$) **apply**(rule $exI[of - w2']$)
apply(rule $exI[of - tr2]$) **apply**(rule $exI[of - s2']$) **apply**(rule $exI[of - s2'']$)
apply(rule $exI[of - ltr2']$)
apply(rule $exI[of - trv1]$) **apply**(rule $exI[of - sv1'']$) **apply**(rule $exI[of - trv2]$)
apply(rule $exI[of - sv2'']$)
apply(rule $exI[of - statOO]$)
apply(intro $conjI$)
subgoal **using** $\chi\chi$.
subgoal **using** $\chi4'$.
subgoal **using** l **unfolding** $h-t.simps$
unfolding $tp-def[symmetric]$ tp **apply** $simp$
unfolding $tp'-def[symmetric]$ tp' **by** auto .
qed

lemma $lltrv1-lltrv2-not-lnever-R$:
assumes unw : $unwindCond \Delta$
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and r : $reachO s1 reachO s2 reachV sv1 reachV sv2$
and $ltr2$: $Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2$
and l' : $\neg isIntO s1 lnever isIntO ltr2$
and len : $llength ltr1 > Suc 0 llength ltr2 > Suc 0$
and l : $trn = R \neg lnever isSecO ltr2$
shows $\exists w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$.
 $\chi\chi s2 ltr2 tr2 s2' s2'' ltr2' \wedge$
 $\chi4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$
 \wedge
 $lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend (l\list-of trv1) (lltrv1 (L, w1', w2', s1, ltr1, s2'', s2'' \$ ltr2', statA, sv1'', sv2'', statOO))$
 \wedge
 $lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend (l\list-of trv2) (lltrv2 (L, w1', w2', s1, ltr1, s2'', s2'' \$ ltr2', statA, sv1'', sv2'', statOO))$

proof –
show ?thesis
using $h-t-not-lnever-R[OF assms]$ **apply**(elim exE)
subgoal **for** $w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$
apply(rule $exI[of - w1']$) **apply**(rule $exI[of - w2']$)
apply(rule $exI[of - tr2]$) **apply**(rule $exI[of - s2']$) **apply**(rule $exI[of - s2'']$)
apply(rule $exI[of - ltr2']$)
apply(rule $exI[of - trv1]$) **apply**(rule $exI[of - sv1'']$) **apply**(rule $exI[of - trv2]$)
apply(rule $exI[of - sv2'']$)
apply(rule $exI[of - statOO]$)

apply(*intro conjI*)
subgoal by simp
subgoal by simp
subgoal unfolding lltrv1-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by simp
subgoal unfolding lltrv2-def apply(subst f-length-ge1[OF len])
unfolding h-def t-def by simp . .
qed

lemma f-not-LNil: $ltr1 \neq [] \vee ltr2 \neq [] \implies$
 $f(w1, w2, trn, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \neq []$
apply(cases llength ltr1 \leq Suc 0 \vee llength ltr2 \leq Suc 0)
subgoal apply(subst f-length-1) **by auto**
subgoal apply(subst f-length-ge1) **by auto .**

lemma lvalidFromS-lltrv1:
assumes unw: *unwindCond* Δ
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and *r:* *reachO* *s1 reachO s2 reachV sv1 reachV sv2*
and *ltr1:* *Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1*
and *ltr2:* *Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2*
shows *Van.lvalidFromS sv1 (lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO))*
proof–
{fix *n sv1 ltrv1*
assume $\exists trn w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO.$
 $n = w1 \wedge$
 $ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$
 $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$
 $Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge lnever isIntO ltr1 \wedge$
 $Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge lnever isIntO ltr2$
hence *Van.llvalidFromS n sv1 ltrv1*
proof(*coinduct rule: Van.llvalidFromS.coinduct[of $\lambda n sv1 ltrv1.$*
 $\exists trn w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO.$
 $n = w1 \wedge$
 $ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$
 $reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$
 $Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge lnever isIntO ltr1 \wedge$
 $Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge lnever isIntO ltr2$])
case (*llvalidFromS n sv1 ltrv1*)
then obtain *trn w1 w2 s1 ltr1 s2 ltr2 statA sv2 statO*
where *n:* $n = w1$ **and**
 $ltrv1: ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$

```

and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and  $r$ :  $reachO s1 reachO s2 reachV sv1 reachV sv2$ 
and  $ltr1$ :  $Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1$ 
and  $ltr2$ :  $Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2$ 
by auto
have  $isi3$ :  $\neg isIntO s1$  using  $ltr1$ 
by (metis  $Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel$ 
 $l\text{list.pred-inject}(2)$ )
have  $isi4$ :  $\neg isIntO s2$  using  $ltr2$ 
by (metis  $Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel$ 
 $l\text{list.pred-inject}(2)$ )

show  $?case$  proof( $cases\ ltr1 = [] \wedge ltr2 = []$ )
case True note  $ltr1_4 = True$ 
hence  $ltrv1$ :  $ltrv1 = []$  unfolding  $ltrv1$  by simp
show  $?thesis$  unfolding  $ltrv1$  apply(rule  $Van.l\text{validFromS-selectLNil}$ ) by
auto
next
case False hence  $ltr1_4$ :  $ltr1 \neq [] \vee ltr2 \neq []$  by auto
show  $?thesis$  proof( $cases\ llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$ )
case True note  $ltr1_4 = ltr1_4\ True$ 
hence  $ltrv1$ :  $ltrv1 = [sv1]$  unfolding  $ltrv1$  by simp
show  $?thesis$  unfolding  $ltrv1$  apply(rule  $Van.l\text{validFromS-selectSingl}$ ) by
auto
next
case False hence  $current$ :  $llength\ ltr1 > Suc\ 0\ llength\ ltr2 > Suc\ 0$ 
by auto
show  $?thesis$  proof( $cases\ trn$ )
case  $L$  note  $trn = L$  note  $current = current\ L$ 
show  $?thesis$ 
proof( $cases\ lnever\ isSecO\ ltr1$ )
case True note  $current = current\ True$ 
obtain  $trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO$  where
 $\omega\omega$ :  $ltr1 = s1 \$ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'$ 
 $lcompletedFromO s1' ltr1' lnever isIntO ltr1'$  and
 $\omega\omega$ :  $\omega\omega\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s1'\ s2\ statA\ sv1\ trv1\ sv1'\ sv2\ trv2\ sv2'$ 
statOO
and  $trn'$ :  $trn' = (if\ trv1 = []\ then\ L\ else\ R)$ 
and  $ltrv1$ :  $ltrv1 =$ 
 $lappend\ (l\text{list-of}\ trv1)\ (lltrv1\ (trn', w1', w2', s1', ltr1', s2, ltr2, statA,$ 
 $sv1', sv2', statOO))$ 
using  $lltrv1-lltrv2-lnever-L[OF\ un\w\ \Delta\ r\ ltr1\ isi4\ current]$ 
unfolding  $ltrv1$  by blast
define  $ltrv1'$  where  $ltrv1'$ :  $ltrv1' \equiv lltrv1\ (trn', w1', w2', s1', ltr1',$ 
 $s2, ltr2, statA, sv1', sv2', statOO)$ 
have  $ltrv1$ :  $ltrv1 = lappend\ (l\text{list-of}\ trv1)\ ltrv1'$ 
unfolding  $ltrv1\ ltrv1' ..$ 

show  $?thesis$ 

```

```

proof(cases trv1 = [])
  case True note trv1 = True
  have sv1': sv1' = sv1
  using  $\omega 3$  unfolding  $\omega 3$ -def by (simp add: trv1)
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w1']) apply(rule exI[of - n])
  apply(rule exI[of - sv1]) apply(rule exI[of - ltrv1'])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1 trv1 by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def trv1 n by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - trn'])
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
    apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
    apply(rule exI[of - statA]) apply(rule exI[of - sv2']) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
    subgoal unfolding ltrv1' trn' trv1 sv1' using trn by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def sv1' by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(1) snd-conv)
    subgoal by fact subgoal by fact
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(4))
    subgoal using  $\omega\omega$  by simp subgoal using  $\omega\omega$  by simp subgoal
using  $\omega\omega$  by simp
    subgoal by fact subgoal by fact subgoal by fact . .
  next
  case False note trv1 = False
  show ?thesis
  apply(rule Van.llvalidFromS-selectlappend)
  apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
  apply(rule exI[of - sv1']) apply(rule exI[of - w1'])
  apply(rule exI[of - ltrv1']) apply(rule exI[of - n])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv1 .
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal by fact
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Van.validFromS-def Van.validS-validTrans list.sel(1)
not-Cons-self2 snoc-eq-iff-butlast trv1)

```

```

      subgoal apply(rule disjI1)
      apply(rule exI[of - trn'])
      apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1']) apply(rule exI[of - ltr1'])
      apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
      apply(rule exI[of - statA]) apply(rule exI[of - sv2']) apply(rule
exI[of - statOO])
      apply(intro conjI)
      subgoal ..
      subgoal using trv1 unfolding ltrv1' trn' by auto
      subgoal using  $\omega$ 3 unfolding  $\omega$ 3-def by simp
      subgoal using  $\omega$ 3 unfolding  $\omega$ 3-def
      by (metis Opt.reach.Step  $\omega$  $\omega$ (2) fst-conv r(1) snd-conv)
      subgoal by fact
      subgoal using  $\omega$ 3 unfolding  $\omega$ 3-def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\omega$ 3 unfolding  $\omega$ 3-def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal using  $\omega$  $\omega$  by auto
      subgoal using  $\omega$  $\omega$  by auto
      subgoal using  $\omega$  $\omega$ 
      using llist-all-lappend-llist-of ltr1(3) by blast
      subgoal using  $\omega$  $\omega$  using ltr2(1) by fastforce
      subgoal by fact
      subgoal by fact ..
    qed
  next
  case False note current = current False
  obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
   $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
   $\chi\mathcal{I}$ ':  $\chi\mathcal{I}$ '  $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
  and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

  using lltrv1-lltrv2-not-lnever-L[OF unw  $\Delta$  r ltr1 isi4 current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

  show ?thesis apply(rule Van.llvalidFromS-selectlappend)
  apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
  apply(rule exI[of - sv1']) apply(rule exI[of - w1'])
  apply(rule exI[of - ltrv1']) apply(rule exI[of - w1])
  apply(intro conjI)

```

```

    subgoal unfolding n .. subgoal ..
    subgoal using ltrv1 .
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by simp
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
by (metis Van.validFromS Van.validS-validTrans Simple-Transition-System.validFromS-def

    append-is-Nil-conv not-Cons-self2)
    subgoal apply (rule disjI1)
    apply (rule exI[of - R])
    apply (rule exI[of - w1 ^]) apply (rule exI[of - w2 ^])
    apply (rule exI[of - s1' ^]) apply (rule exI[of - s1'' $ ltr1 ^])
    apply (rule exI[of - s2]) apply (rule exI[of - ltr2])
    apply (rule exI[of - statA]) apply (rule exI[of - sv2' ^]) apply (rule exI[of
- statOO])
    apply (intro conjI)
    subgoal ..
    subgoal unfolding ltrv1' ..
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by simp
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

    append-is-Nil-conv last-snoc not-Cons-self2 r(1))
    subgoal by fact
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr1(3) by blast
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
    subgoal by fact
    subgoal by fact . .

qed
next
case R note trn = R note current = current R
show ?thesis
proof (cases lnever isSecO ltr2)
case True note current = current True
obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
 $\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
 $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 

```

```

statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv1: ltrv1 =
    lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF unu Δ r ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

show ?thesis
proof(cases trv1 = [])
  case True note trv1 = True
  have sv1': sv1' = sv1
  using ω4 unfolding ω4-def by (simp add: trv1)
  show ?thesis
  apply(rule Van.llvalidFromS-selectDelay)
  apply(rule exI[of - w1]) apply(rule exI[of - n])
  apply(rule exI[of - sv1]) apply(rule exI[of - ltrv1])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1 trv1 by simp
  subgoal using ω4 unfolding ω4-def trv1 n by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - trn])
    apply(rule exI[of - w1]) apply(rule exI[of - w2])
    apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv2]) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
    subgoal unfolding ltrv1' trn' trv1 sv1' using trn by simp
    subgoal using ω4 unfolding ω4-def sv1' by simp
    subgoal by fact
    subgoal using ω4 unfolding ω4-def
    by (metis Opt.reach.Step ωω(2) fst-conv r(2) snd-conv)
    subgoal by fact
    subgoal using ω4 unfolding ω4-def
    by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using ωω by simp subgoal using ωω by simp subgoal
using ωω by simp . .
  next
  case False note trv1 = False
  show ?thesis

```

```

apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1']) apply(rule exI[of - w1'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv1 .
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Van.validFromS-def Van.validS-validTrans list.sel(1)
not-Cons-self2 snoc-eq-iff-butlast trv1)
  subgoal apply(rule disjI1)
  apply(rule exI[of - trn'])
  apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
  apply(rule exI[of - statA]) apply(rule exI[of - sv2']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal using trv1 unfolding ltrv1' trn' by auto
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast . .
qed
next
case False note current = current False
obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
 $\chi\chi$ :  $\chi\chi$  s2 ltr2 tr2 s2' s2'' ltr2' and
 $\chi\chi'$ :  $\chi\chi'$   $\Delta$  w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))

```



```

    using lltrv1-lltrv2-not-lnever-R[OF un $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
  have ltrv1: ltrv1 = lappend (l $\text{list-of}$  trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

  show ?thesis
  proof (cases trv1 = [])
    case True note trv1 = True
    hence sv1'': sv1'' = sv1
    by (metis  $\chi_4'$ -def Simple-Transition-System.validFromS-Cons-iff  $\chi_4'$ 
append.simps(1))
    have w1' < w1 using trv1  $\chi_4'$  unfolding  $\chi_4'$ -def by auto
    show ?thesis
    apply (rule Van.llvalidFromS-selectDelay)
    apply (rule exI[of - w1']) apply (rule exI[of - n])
    apply (rule exI[of - sv1]) apply (rule exI[of - ltrv1])
    apply (intro conjI)
    subgoal ..
    subgoal .. subgoal .. subgoal unfolding n by fact
    subgoal apply (rule disjI1)
    apply (rule exI[of - L])
    apply (rule exI[of - w1']) apply (rule exI[of - w2'])
    apply (rule exI[of - s1]) apply (rule exI[of - ltr1])
    apply (rule exI[of - s2'']) apply (rule exI[of - s2'' $ ltr2'])
    apply (rule exI[of - statA]) apply (rule exI[of - sv2'']) apply (rule
exI[of - statOO])
    apply (intro conjI)
    subgoal ..
    subgoal unfolding ltrv1 ltrv1' trv1 sv1'' by simp
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def sv1'' by simp
    subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    by (metis Opt.reach-validFromS-reach Nil-is-append-conv last-snoc
not-Cons-self2 r(2))
    subgoal by fact
    subgoal using  $\chi_4'$  r(4) unfolding  $\chi_4'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
    using llist-all-lappend-llist-of ltr2(3) by blast ..
  next
  case False note trv1 = False
  show ?thesis
  apply (rule Van.llvalidFromS-selectlappend)

```

```

apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1'])
apply(rule exI[of - w1'])
apply(rule exI[of - ltrv1'])
apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv1 .
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal using trv1 .
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
list.sel(1))
    not-Cons-self2 snoc-eq-iff-butlast trv1)
  subgoal apply(rule disjI1)
  apply(rule exI[of - L])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
  apply(rule exI[of - s2'']) apply(rule exI[of - s2''] $ ltr2')
  apply(rule exI[of - statA]) apply(rule exI[of - sv2'']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal unfolding ltrv1' ..
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
  subgoal by fact
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$ 
xx-def)
  append-is-Nil-conv last-snoc not-Cons-self2 r(2))
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal by fact
  subgoal by fact
  subgoal by fact
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  using l1ist-all-lappend-l1ist-of ltr2(3) by blast . .
  qed
  qed
  qed
  qed
  qed

```

```

qed
}
thus ?thesis apply-apply(rule Van.lvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

```

lemma lvalidFromS-lltrv2:
assumes unw: unwindCond Δ
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
shows Van.lvalidFromS sv2 (lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof -
  {fix n sv2 ltrv2
  assume ∃ trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO.
    n = w2 ∧
    ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧

    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2
  hence Van.lvalidFromS n sv2 ltrv2
  proof (coinduct rule: Van.lvalidFromS.coinduct[of λn sv2 ltrv2.
    ∃ trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO.
      n = w2 ∧
      ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
      Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
      reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
      Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧

      Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2])
  case (llvalidFromS n sv2 ltrv2)
  then obtain trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO
  where n: n = w2 and
  ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  by auto
  have isi3: ¬ isIntO s1 using ltr1
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
  llist.pred-inject(2))
  have isi4: ¬ isIntO s2 using ltr2
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
  llist.pred-inject(2))
  }

```

```

show ?case proof(cases  $ltr1 = [] \wedge ltr2 = []$ )
  case True note  $ltr1_4 = True$ 
  hence  $ltrv2: ltrv2 = []$  unfolding  $ltrv2$  by simp
  show ?thesis unfolding  $ltrv2$  apply(rule Van.lvalidFromS-selectLNil) by
auto
  next
  case False hence  $ltr1_4: ltr1 \neq [] \vee ltr2 \neq []$  by auto
  show ?thesis proof(cases  $l\text{length } ltr1 \leq Suc\ 0 \vee l\text{length } ltr2 \leq Suc\ 0$ )
    case True note  $ltr1_4 = ltr1_4\ True$ 
    hence  $ltrv2: ltrv2 = [[sv2]]$  unfolding  $ltrv2$  by simp
    show ?thesis unfolding  $ltrv2$  apply(rule Van.lvalidFromS-selectSingl) by
auto
    next
    case False hence current:  $l\text{length } ltr1 > Suc\ 0\ l\text{length } ltr2 > Suc\ 0$ 
    by auto
    show ?thesis proof(cases  $trn$ )
      case L note  $trn = L$  note current = current L
      show ?thesis
      proof(cases  $l\text{never } isSecO\ ltr1$ )
        case True note current = current True
        obtain  $trn'\ w1'\ w2'\ s1'\ ltr1'\ trv1\ sv1'\ trv2\ sv2'\ statOO$  where
         $\omega\omega: ltr1 = s1\ \$\ ltr1'\ \text{validTransO}\ (s1, s1')\ \text{Opt.lvalidFromS}\ s1'\ ltr1'$ 
         $l\text{completedFromO}\ s1'\ ltr1'\ l\text{never } isIntO\ ltr1'$  and
         $\omega3: \omega3\ \Delta\ w1\ w2\ w1'\ w2'\ s1\ s1'\ s2\ statA\ sv1\ trv1\ sv1'\ sv2\ trv2\ sv2'$ 
statOO
        and  $trn': trn' = (\text{if } trv1 = [] \text{ then } L \text{ else } R)$ 
        and  $ltrv2: ltrv2 =$ 
         $l\text{append}\ (l\text{list-of } trv2)\ (l\text{ltrv2}(trn', w1', w2', s1', ltr1', s2, ltr2, statA,$ 
sv1', sv2', statOO))
        using  $lltrv1\ lltrv2\ l\text{never}\ L[OF\ un\ \Delta\ r\ ltr1\ isi4\ \text{current}]$ 
        unfolding  $ltrv2$  by blast
        define  $ltrv2'$  where  $ltrv2': ltrv2' \equiv l\text{ltrv2}\ (trn', w1', w2', s1', ltr1',$ 
s2, ltr2, statA, sv1', sv2', statOO)
        have  $ltrv2: ltrv2 = l\text{append}\ (l\text{list-of } trv2)\ ltrv2'$ 
        unfolding  $ltrv2\ ltrv2' ..$ 

        show ?thesis
        proof(cases  $trv2 = []$ )
          case True note  $trv2 = True$ 
          have  $sv2': sv2' = sv2$ 
          using  $\omega3$  unfolding  $\omega3\text{-def}$  by (simp add: trv2)
          show ?thesis
          apply(rule Van.lvalidFromS-selectDelay)
          apply(rule  $exI[of - w2']$ ) apply(rule  $exI[of - n]$ )
          apply(rule  $exI[of - sv2]$ ) apply(rule  $exI[of - ltrv2']$ )
          apply(intro conjI)
          subgoal .. subgoal ..
          subgoal unfolding  $ltrv2\ trv2$  by simp

```

```

subgoal using  $\omega 3$  unfolding  $\omega 3$ -def trv2 n by simp
subgoal apply(rule disjI1)
  apply(rule exI[of - trn'])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
    subgoal unfolding ltrv2' trn' trv2 sv2' using trn by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def sv2' by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(3))
    subgoal by fact
    subgoal using  $\omega\omega$  by simp subgoal using  $\omega\omega$  by simp subgoal
using  $\omega\omega$  by simp
    subgoal by fact subgoal by fact subgoal by fact . .
  next
  case False note trv2 = False
  show ?thesis
  apply(rule Van.llvalidFromS-selectlappend)
  apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
  apply(rule exI[of - sv2']) apply(rule exI[of - w2'])
  apply(rule exI[of - ltrv2']) apply(rule exI[of - n])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv2 .
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal by fact
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Van.validFromS-def Van.validS-validTrans append-is-Nil-conv
list.sel(1) not-Cons-self2 trv2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - trn'])
  apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1']) apply(rule exI[of - ltr1'])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal using trv2 unfolding ltrv2' trn' by auto

```

```

    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
    subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
    by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$ 
    using llist-all-lappend-llist-of ltr1(3) by blast
    subgoal using  $\omega\omega$  using ltr2(1) by fastforce
    subgoal by fact
    subgoal by fact . .
qed
next
case False note current = current False
obtain  $w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO$  where
 $\chi\chi: \chi\chi s1 ltr1 tr1 s1' s1'' ltr1'$  and
 $\chi 3': \chi 3' \Delta w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2$ 
 $trv2 sv2'' statOO$ 
and  $ltrv2: ltrv2 =$ 
lappend (llist-of trv2) (lltrv2 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

using lltrv1-lltrv2-not-lnever-L[OF unw  $\Delta r ltr1 isi4 current$ ]
unfolding ltrv2 by blast
define ltrv2' where  $ltrv2': ltrv2' \equiv lltrv2 (R,w1',w2',s1'',s1'' $$ 
 $ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)$ 
have  $ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'$ 
unfolding ltrv2 ltrv2' ..

show ?thesis
proof(cases  $trv2 = []$ )
case True note  $trv2 = True$ 
hence  $sv2'': sv2'' = sv2$ 
by (metis  $\chi 3'$ -def Simple-Transition-System.validFromS-Cons-iff  $\chi 3'$ 
append.simps(1))
have  $w2' < w2$  using  $trv2 \chi 3'$  unfolding  $\chi 3'$ -def by auto
show ?thesis
apply(rule Van.llvalidFromS-selectDelay)
apply(rule exI[of -  $w2'$ ]) apply(rule exI[of -  $n$ ])
apply(rule exI[of -  $sv2$ ]) apply(rule exI[of -  $ltrv2$ ])
apply(intro conjI)
subgoal ..
subgoal .. subgoal .. subgoal unfolding  $n$  by fact
subgoal apply(rule disjI1)

```

```

apply(rule exI[of - R])
apply(rule exI[of - w1  $\uparrow$ ]) apply(rule exI[of - w2  $\uparrow$ ])
apply(rule exI[of - s1  $\uparrow$ ]) apply(rule exI[of - s1'' $ ltr1  $\uparrow$ ])
apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1  $\uparrow$ ]) apply(rule
exI[of - statOO])
apply(intro conjI)
  subgoal ..
    subgoal unfolding ltrv2 ltrv2' trv2 sv2'' by simp
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  sv2'' by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
    by (metis Opt.reach-validFromS-reach Nil-is-append-conv last-snoc
not-Cons-self2 r(1))
    subgoal by fact
    subgoal using  $\chi^{3'}$  r(3) unfolding  $\chi^{3'-def}$ 
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
      using llist-all-lappend-llist-of ltr1(3) by blast
    subgoal by fact subgoal by fact subgoal by fact . .
  next
  case False note trv2 = False
  show ?thesis
  apply(rule Van.lvalidFromS-selectlappend)
  apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
  apply(rule exI[of - sv2  $\uparrow$ ]) apply(rule exI[of - w2  $\uparrow$ ])
  apply(rule exI[of - ltrv2  $\uparrow$ ]) apply(rule exI[of - w2])
  apply(intro conjI)
    subgoal unfolding n .. subgoal ..
    subgoal using ltrv2 .
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
    subgoal by fact
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
append-is-Nil-conv list.sel(1) not-Cons-self2 trv2)
    subgoal apply(rule disjI1)
    apply(rule exI[of - R])
    apply(rule exI[of - w1  $\uparrow$ ]) apply(rule exI[of - w2  $\uparrow$ ])
    apply(rule exI[of - s1  $\uparrow$ ]) apply(rule exI[of - s1'' $ ltr1  $\uparrow$ ])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
    apply(rule exI[of - statA]) apply(rule exI[of - sv1  $\uparrow$ ]) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
    subgoal unfolding ltrv2' ..

```

```

      subgoal using  $\chi^3'$  unfolding  $\chi^3'$ -def by simp
      subgoal using  $\chi^3'$  unfolding  $\chi^3'$ -def
        by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$ 
 $\chi\chi$ -def
          append-is-Nil-conv last-snoc not-Cons-self2 r(1))
      subgoal by fact
      subgoal using  $\chi^3'$  unfolding  $\chi^3'$ -def
        by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\chi^3'$  unfolding  $\chi^3'$ -def
        by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
        using llist-all-lappend-llist-of ltr1(3) by blast
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
      subgoal by fact
      subgoal by fact . .
    qed
  qed
next
case R note trn = R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
 $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
 $\omega_4$ :  $\omega_4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv2: ltrv2 =
lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF un $\omega \Delta r$  ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..

  show ?thesis
  proof(cases trv2 = [])
    case True note trv2 = True
    have sv2': sv2' = sv2
    using  $\omega_4$  unfolding  $\omega_4$ -def by (simp add: trv2)
    show ?thesis
    apply(rule Van.llvalidFromS-selectDelay)

```



```

apply(rule exI[of - w2']) apply(rule exI[of - n])
apply(rule exI[of - sv2]) apply(rule exI[of - ltrv2'])
apply(intro conjI)
subgoal .. subgoal ..
subgoal unfolding ltrv2 trv2 by simp
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def trv2 n by simp
subgoal apply(rule disjI1)
  apply(rule exI[of - trn'])
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
  apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding ltrv2' trn' trv2 sv2' using trn by simp
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def sv2' by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Nil-is-append-conv Van.reach-validFromS-reach last-snoc
not-Cons-self2 r(3))
  subgoal by fact subgoal by fact subgoal by fact subgoal by
fact
  subgoal using  $\omega\omega$  by simp subgoal using  $\omega\omega$  by simp subgoal
using  $\omega\omega$  by simp . .
next
case False note trv2 = False
show ?thesis
apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
apply(rule exI[of - sv2']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv2']) apply(rule exI[of - n])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv2 .
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Van.validFromS-def Van.validS-validTrans append-is-Nil-conv
list.sel(1) not-Cons-self2 trv2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - trn'])
  apply(rule exI[of - w1']) apply(rule exI[of - w2']) apply(rule exI[of
- s1]) apply(rule exI[of - ltr1])
  apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])

```

```

      apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule
exI[of - statOO])
    apply(intro conjI)
    subgoal ..
      subgoal using trv2 unfolding ltrv2' trn' by auto
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
      subgoal by fact
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
    by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$ 
      using llist-all-lappend-llist-of ltr1(3) by blast . .
    qed
  next
  case False note current = current False
  obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
   $\chi\chi$ :  $\chi\chi$  s2 ltr2 tr2 s2' s2'' ltr2' and
   $\chi_4'$ :  $\chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
trv2 sv2'' statOO
  and ltrv2: ltrv2 =
    lappend (llist-of trv2) (lltrv2 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
  have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..
  have trv2: trv2  $\neq$  [] using  $\chi_4'$  unfolding  $\chi_4'$ -def by auto

  show ?thesis
  apply(rule Van.lvalidFromS-selectlappend)
  apply(rule exI[of - sv2]) apply(rule exI[of - trv2])
  apply(rule exI[of - sv2'])
  apply(rule exI[of - w2'])
  apply(rule exI[of - ltrv2'])
  apply(rule exI[of - n])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal using ltrv2 .

```

```

      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Nil-is-append-conv Van.validFromS-def Van.validS-append1
hd-append2)
      subgoal using trv2 .
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
append-is-Nil-conv list.sel(1) not-Cons-self2)
      subgoal apply(rule disjI1)
      apply(rule exI[of - L])
      apply(rule exI[of - w1']) apply(rule exI[of - w2'])
      apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
      apply(rule exI[of - s2']) apply(rule exI[of - s2'' $ ltr2'])
      apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule exI[of
- statOO])
      apply(intro conjI)
      subgoal .. subgoal unfolding ltrv2' ..
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
      subgoal by fact
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def
append-is-Nil-conv last-snoc not-Cons-self2 r(2))
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal by fact
      subgoal by fact
      subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast . .
      qed
    qed
  qed
}
}
thus ?thesis apply-apply(rule Van.lvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

lemma lcompletedFrom-lltrv1:
 assumes unw: unwindCond Δ

```

and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and  $r$ :  $reachO s1 reachO s2 reachV sv1 reachV sv2$ 
and  $ltr1$ :  $Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1$ 
and  $ltr2$ :  $Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2$ 
shows  $Van.lcompletedFrom sv1 (lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))$ 
proof –
  {fix  $ltrv1$  assume  $ltrv1$ :  $ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)$ 
    and  $lfin$ :  $lfinite ltrv1$ 
    hence  $list-of ltrv1 \neq [] \wedge finalV (last (list-of ltrv1))$ 
    using  $assms(2-)$  proof( $induct\ length\ (list-of\ ltrv1)\ w1$ 
       $arbitrary$ :  $trn\ w2\ ltrv1\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ sv2\ statO$ 
       $rule$ :  $less2-induct'$ )
    case ( $less\ w1\ ltrv1\ trn\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv1\ sv2\ statO$ )
      hence  $ltrv1$ :  $ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,$ 
 $statO)$ 
      and  $lfin$ :  $lfinite ltrv1$ 
      and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
      and  $r$ :  $reachO s1 reachO s2 reachV sv1 reachV sv2$ 
      and  $ltr1$ :  $Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1 lnever isIntO ltr1$ 
      and  $ltr2$ :  $Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2 lnever isIntO ltr2$ 
      by  $auto$ 
      have  $isi3$ :  $\neg isIntO s1$  using  $ltr1$ 
      by ( $metis\ Opt.lcompletedFrom-def\ Opt.lvalidFromS-def\ lfinite-LNil\ llist.exhaust-sel$ 
 $l\ list.pred-inject(2)$ )
      have  $isi4$ :  $\neg isIntO s2$  using  $ltr2$ 
      by ( $metis\ Opt.lcompletedFrom-def\ Opt.lvalidFromS-def\ lfinite-LNil\ llist.exhaust-sel$ 
 $l\ list.pred-inject(2)$ )

      show  $?case$  proof( $cases\ ltr1 = [[]] \wedge ltr2 = [[]]$ )
        case  $True$  note  $ltr14 = True$ 
        hence  $False$  using  $ltr1(2)\ ltr2(2)$  unfolding  $Opt.lcompletedFrom-def$  by
 $auto$ 

        thus  $?thesis$  by  $auto$ 
      next
      case  $False$  hence  $ltr14$ :  $ltr1 \neq [[]] \vee ltr2 \neq [[]]$  by  $auto$ 
      show  $?thesis$  proof( $cases\ llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0$ )
        case  $True$  note  $ltr14 = ltr14\ True$ 
        hence  $ltrv1$ :  $list-of\ ltrv1 = [sv1]$  unfolding  $ltrv1$  by  $simp$ 
        have  $llength\ ltr1 = Suc\ 0 \vee llength\ ltr2 = Suc\ 0$ 
        using  $ltr14$ 
        by ( $metis\ Opt.lcompletedFrom-def$ 
           $Suc-ile-eq\ i0-less\ lfinite-code(1)\ llength-eq-0\ llist.exhaust$ 
 $ltr1(2)\ ltr2(2)\ nle-le\ not-lnull-conv\ zero-enat-def$ )
        hence  $ltr1 = [[s1]] \vee ltr2 = [[s2]]$ 
        using  $Opt.lcompletedFrom-singl\ ltr1(1)\ ltr1(2)\ ltr2(1)\ ltr2(2)$  by  $blast$ 
        hence  $finalO\ s1 \vee finalO\ s2$ 
        using  $Opt.lcompletedFrom-LCons\ ltr1(2)\ ltr2(2)$  by  $blast$ 
        hence  $finalV\ sv1$ 
        using  $\Delta\ r(1)\ r(2)\ r(3)\ r(4)\ unw\ unwindCond-def$  by  $auto$ 

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thus ?thesis unfolding ltrv1 by auto
next
  case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0 by
auto
show ?thesis
proof(cases trn)
  case L note current = current L
  show ?thesis
  proof(cases lnever isSecO ltr1)
    case True note current = current True
    obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
    ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
    lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
    ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
    and trn' : trn' = (if trv1 = [] then L else R)
    and lltrv1: ltrv1 =
    lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unω Δ r ltr1 isi4 current]
    unfolding ltrv1 by blast
    define ltrv1' where ltrv1': ltrv1' = lltrv1 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
    have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
    unfolding lltrv1 ltrv1' ..

    have trv1ne: trv1 ≠ [] ∨ w1' < w1 using ω3 unfolding ω3-def by
auto
    have lfin': lfinite ltrv1'
    using lfin trv1ne unfolding lltrv1 by simp
    have len: length (list-of ltrv1') < length (list-of ltrv1) ∨
    length (list-of ltrv1') = length (list-of ltrv1) ∧ w1' < w1
    using trv1ne lfin lfin' by (simp add: list-of-lappend lltrv1)

    have 0: list-of ltrv1' ≠ [] ∧ finalV (last (list-of ltrv1'))
    using len proof(elim disjE conjE)
    assume len: length (list-of ltrv1') < length (list-of ltrv1)
    show ?thesis
    apply(rule less(1)[OF - ltrv1'])
    subgoal by fact subgoal by fact
    subgoal using ω3 unfolding ω3-def by simp
    subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact

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```

      subgoal by fact subgoal by fact .
next
assume len: length (list-of ltrv1') = length (list-of ltrv1) w1' < w1
show ?thesis
apply(rule less(2)[OF - - ltrv1'])
      subgoal by fact subgoal using len by simp subgoal by fact

      subgoal using ω3 unfolding ω3-def by simp
      subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
      subgoal by fact
      subgoal using ω3 unfolding ω3-def
      by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
      subgoal using ω3 unfolding ω3-def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
      subgoal by fact subgoal by fact subgoal by fact subgoal by fact
      subgoal by fact subgoal by fact .
qed
show ?thesis unfolding lltrv1 using 0
by (simp add: lfin' list-of-lappend)
next
case False note current = current False
obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
χχ: χχ s1 ltr1 tr1 s1' s1'' ltr1' and
χ3': χ3' Δ w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and lltrv1: ltrv1 =
lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

using lltrv1-lltrv2-not-lnever-L[OF unω Δ r ltr1 isi4 current]
unfolding ltrv1 by blast
define ltrv1' where ltrv1': ltrv1' = lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)

have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
unfolding lltrv1 ltrv1' ..
have trv1ne: trv1 ≠ [] using χ3' unfolding χ3'-def by auto
have lfin': lfinite ltrv1'
using lfin trv1ne unfolding lltrv1 by simp
have len: length (list-of ltrv1') < length (list-of ltrv1)
using trv1ne lfin lfin' by (simp add: list-of-lappend lltrv1)

have 0: list-of ltrv1' ≠ [] ∧ finalV (last (list-of ltrv1'))
apply(rule less(1)[OF - ltrv1'])
subgoal by fact subgoal by fact
subgoal using χ3' unfolding χ3'-def by simp
subgoal using χχ unfolding χχ-def
by (metis Simple-Transition-System.reach-validFromS-reach r(1)
snoc-eq-iff-butlast)

```

```

      subgoal by fact
      subgoal using  $\chi 3'$  unfolding  $\chi 3'$ -def
        by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(3))
      subgoal using  $\chi 3'$  unfolding  $\chi 3'$ -def
        by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
      subgoal using  $\chi X$  unfolding  $\chi X$ -def by simp
      subgoal using  $\chi X$  unfolding  $\chi X$ -def by simp
      subgoal using  $\chi X$  unfolding  $\chi X$ -def
        using llist-all-lappend-llist-of ltr1 by blast
      subgoal by fact subgoal by fact subgoal by fact .
      show ?thesis unfolding lltrv1 using 0
        by (simp add: lfin' list-of-lappend)
    qed
  next
  case R note current = current R
  show ?thesis
  proof(cases lnever isSecO ltr2)
    case True note current = current True
    obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
       $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
      lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
       $\omega 4$ :  $\omega 4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 
statOO
    and trn': trn' = (if trv2 = [] then R else L)
    and ltrv1: ltrv1 =
      lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-R[OF unW  $\Delta$  r ltr2(1,2) isi3 ltr2(3) current]
    unfolding ltrv1 by blast
    define ltrv1' where ltrv1': ltrv1' = lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
    have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
    unfolding ltrv1 ltrv1' ..

    have trv1ne: trv1  $\neq$  []  $\vee$  w1' < w1 using  $\omega 4$  unfolding  $\omega 4$ -def by
auto
    have lfin': lfinite ltrv1'
    using lfin trv1ne unfolding lltrv1 by simp
    have len: length (list-of ltrv1') < length (list-of ltrv1)  $\vee$ 
      length (list-of ltrv1') = length (list-of ltrv1)  $\wedge$  w1' < w1
    using trv1ne lfin lfin' by (simp add: list-of-lappend lltrv1)

    have 0: list-of ltrv1'  $\neq$  []  $\wedge$  finalV (last (list-of ltrv1'))
    using len proof(elim disjE conjE)
      assume len: length (list-of ltrv1') < length (list-of ltrv1)
      show ?thesis
      apply(rule less(1)[OF - ltrv1'])

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```

    subgoal by fact subgoal by fact
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
    subgoal by fact
    subgoal using  $r(2)$   $\omega\omega$  by (metis Opt.reach.Step fst-conv snd-conv)
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2  $r(3)$ )
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2  $r(4)$ )
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact
    subgoal by fact subgoal by fact .
  next
  assume len: length (list-of ltrv1') = length (list-of ltrv1)  $w1' < w1$ 
  show ?thesis
  apply (rule less(2)[OF - - ltrv1'])
    subgoal by fact subgoal using len by simp subgoal by fact

    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
    subgoal by fact
    subgoal by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv  $r(2)$  snd-conv)
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach  $r(3)$  snoc-eq-iff-butlast)
    subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2  $r(4)$ )
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact
    subgoal by fact subgoal by fact .
  qed
  show ?thesis unfolding lltrv1 using 0
  by (simp add: lfn' list-of-lappend)
next
case False note current = current False
obtain  $w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$  where
 $\chi\chi: \chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$  and
 $\chi_4': \chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
 $trv2 sv2'' statOO$ 
and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (L,  $w1'$ ,  $w2'$ ,  $s1$ , ltr1,  $s2''$ ,  $s2''$  $ ltr2',
statA,  $sv1''$ ,  $sv2''$ , statOO))
  using lltrv1-lltrv2-not-lnever-R[OF un $w$   $\Delta$   $r$  ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' = lltrv1 (L,  $w1'$ ,  $w2'$ ,  $s1$ , ltr1,  $s2''$ ,
 $s2''$  $ ltr2', statA,  $sv1''$ ,  $sv2''$ , statOO)
  have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

  have trv1ne: trv1  $\neq$  []  $\vee w1' < w1$  using  $\chi_4'$  unfolding  $\chi_4'$ -def by

```


auto

```
have lfin': lfinite ltrv1'
using lfin trv1ne unfolding ltrv1 by simp
have len: length (list-of ltrv1') < length (list-of ltrv1) ∨
      length (list-of ltrv1') = length (list-of ltrv1) ∧ w1' < w1
using trv1ne lfin lfin' by (simp add: list-of-lappend ltrv1)

have 0: list-of ltrv1' ≠ [] ∧ finalV (last (list-of ltrv1'))
using len proof (elim disjE conjE)
  assume len: length (list-of ltrv1') < length (list-of ltrv1)
  show ?thesis
  apply (rule less(1)[OF - ltrv1'])
  subgoal by fact subgoal by fact
  subgoal using χ4' unfolding χ4'-def by simp
  subgoal by fact
  subgoal using r(2) χχ unfolding χχ-def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using χ4' unfolding χ4'-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(3))
  subgoal using χ4' unfolding χ4'-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using χχ unfolding χχ-def by auto
  subgoal using χχ unfolding χχ-def by auto
  subgoal using χχ unfolding χχ-def
    using llist-all-lappend-llist-of ltr2(3) by blast .
next
assume len: length (list-of ltrv1') = length (list-of ltrv1) w1' < w1
show ?thesis
apply (rule less(2)[OF - - ltrv1'])
  subgoal by fact subgoal using len by simp subgoal by fact

  subgoal using χ4' unfolding χ4'-def by simp
  subgoal by fact
  subgoal using χχ unfolding χχ-def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(2))
  subgoal using χ4' unfolding χ4'-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using χ4' unfolding χ4'-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(4))

  subgoal by fact subgoal by fact subgoal by fact
  subgoal using χχ unfolding χχ-def by auto
  subgoal using χχ unfolding χχ-def by auto
  subgoal using χχ unfolding χχ-def
```

```

        using llist-all-lappend-llist-of-ltr2(3) by blast .
      qed
      show ?thesis unfolding lltrv1 using 0
      by (simp add: lfin' list-of-lappend)
    qed
  qed
  qed
  qed
  qed
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

lemma lcompletedFrom-lltrv2:
  assumes unw: unwindCond  $\Delta$ 
  and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  shows Van.lcompletedFrom sv2 (lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  proof -
    {fix ltrv2 assume ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
      and lfin: lfinite ltrv2
      hence list-of ltrv2  $\neq [] \wedge$  finalV (last (list-of ltrv2))
      using assms(2-) proof(induct length (list-of ltrv2) w2
        arbitrary: ltrv2 trn w1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
        rule: less2-induct')
        case (less w2 ltrv2 trn w1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)
          hence ltrv2: ltrv2 = lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
            statO)
          and lfin: lfinite ltrv2
          and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
          and r: reachO s1 reachO s2 reachV sv1 reachV sv2
          and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1 lnever isIntO ltr1
          and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2 lnever isIntO ltr2
          by auto
          have isi3:  $\neg$  isIntO s1 using ltr1
          by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
            llist.pred-inject(2))
          have isi4:  $\neg$  isIntO s2 using ltr2
          by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
            llist.pred-inject(2))

          show ?case proof(cases ltr1 = [[]]  $\wedge$  ltr2 = [[]])
            case True note ltr14 = True
            hence False using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by
              auto
            thus ?thesis by auto
          next

```

```

case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
  hence ltrv2: list-of ltrv2 = [sv2] unfolding ltrv2 by simp
  have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
  using ltr14
  by (metis Opt.lcompletedFrom-def
    Suc-ile-eq i0-less lfinite-code(1) llength-eq-0 llist.exhaust
    ltr1(2) ltr2(2) nle-le not-lnull-conv zero-enat-def)
  hence ltr1 = [[s1]] ∨ ltr2 = [[s2]]
  using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
  hence finalO s1 ∨ finalO s2
  using Opt.lcompletedFrom-LCons ltr1(2) ltr2(2) by blast
  hence finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unw unwindCond-def by auto
  thus ?thesis unfolding ltrv2 by auto
next
  case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0 by
auto
  show ?thesis
  proof(cases trn)
    case L note current = current L
    show ?thesis
    proof(cases lnever isSecO ltr1)
      case True note current = current True
      obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
         $\omega\omega$ : ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
         $\omega\omega\omega$ :  $\omega\omega\omega$   $\Delta$  w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
        and trn' : trn' = (if trv1 = [] then L else R)
        and lltrv2: ltrv2 =
          lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
        using lltrv1-lltrv2-lnever-L[OF unw  $\Delta$  r ltr1 isi4 current]
        unfolding ltrv2 by blast
        define ltrv2' where ltrv2': ltrv2' = lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
        have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
        unfolding lltrv2 ltrv2' ..

        have trv2ne: trv2 ≠ [] ∨ w2' < w2 using  $\omega\omega\omega$  unfolding  $\omega\omega\omega$ -def by
auto
        have lfin': lfinite ltrv2'
        using lfin trv2ne unfolding lltrv2 by simp
        have len: length (list-of ltrv2') < length (list-of ltrv2) ∨
          length (list-of ltrv2') = length (list-of ltrv2)  $\wedge$  w2' < w2
        using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

```

```

have 0: list-of ltrv2' ≠ [] ∧ finalV (last (list-of ltrv2'))
using len proof(elim disjE conjE)
  assume len: length (list-of ltrv2') < length (list-of ltrv2)
  show ?thesis
  apply(rule less(1)[OF - ltrv2'])
    subgoal by fact subgoal by fact
    subgoal using ω3 unfolding ω3-def by simp
    subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact
    subgoal by fact subgoal by fact .
  next
  assume len: length (list-of ltrv2') = length (list-of ltrv2) w2' < w2
  show ?thesis
  apply(rule less(2)[OF - - ltrv2'])
    subgoal by fact subgoal using len by simp subgoal by fact

    subgoal using ω3 unfolding ω3-def by simp
    subgoal by (metis Opt.reach.Step ωω(2) fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using ω3 unfolding ω3-def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal by fact subgoal by fact subgoal by fact subgoal by fact
    subgoal by fact subgoal by fact .
  qed
  show ?thesis unfolding lltrv2 using 0
  by (simp add: lfin' list-of-lappend)
next
  case False note current = current False
  obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
  χχ: χχ s1 ltr1 tr1 s1' s1'' ltr1' and
  χ3': χ3' Δ w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
  and lltrv2: ltrv2 =
lappend (llist-of trv2) (lltrv2 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

  using lltrv1-lltrv2-not-lnever-L[OF unω Δ r ltr1 isi4 current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2' = lltrv2 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'

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unfolding lltrv2 ltrv2' ..

have trv2ne: trv2 ≠ [] ∨ w2' < w2 using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by
auto

have lfin': lfinite ltrv2'
using lfin trv2ne unfolding lltrv2 by simp
have len: length (list-of ltrv2') < length (list-of ltrv2) ∨
          length (list-of ltrv2') = length (list-of ltrv2) ∧ w2' < w2
using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

have 0: list-of ltrv2' ≠ [] ∧ finalV (last (list-of ltrv2'))
using len proof(elim disjE conjE)
  assume len: length (list-of ltrv2') < length (list-of ltrv2)
  show ?thesis
  apply(rule less(1)[OF - ltrv2'])
    subgoal by fact subgoal by fact
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
      by (metis Simple-Transition-System.reach-validFromS-reach r(1))
snoc-eq-iff-butlast)
    subgoal by fact
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
r(3))
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
      using llist-all-lappend-llist-of ltr1 by blast
    subgoal by fact subgoal by fact subgoal by fact .
next
  assume len: length (list-of ltrv2') = length (list-of ltrv2) w2' < w2
  show ?thesis
  apply(rule less(2)[OF - - ltrv2'])
    subgoal by fact subgoal using len by simp subgoal by fact

    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by simp
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(1))
    subgoal by fact
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
      by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
    subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto

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```

      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
        using llist-all-lappend-llist-of ltr1(3) by blast
      subgoal by fact subgoal by fact subgoal by fact .
    qed
    show ?thesis unfolding lltrv2 using 0
    by (simp add: lfin' list-of-lappend)
  qed
next
case R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
   $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
  lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
   $\omega4$ :  $\omega4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF un $\omega \Delta r$  ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2' = lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..

  have trv2ne: trv2  $\neq$  []  $\vee$  w2' < w2 using  $\omega4$  unfolding  $\omega4$ -def by
auto
  have lfin': lfinite ltrv2'
  using lfin trv2ne unfolding lltrv2 by simp
  have len: length (list-of ltrv2') < length (list-of ltrv2)  $\vee$ 
  length (list-of ltrv2') = length (list-of ltrv2)  $\wedge$  w2' < w2
  using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

  have 0: list-of ltrv2'  $\neq$  []  $\wedge$  finalV (last (list-of ltrv2'))
  using len proof(elim disjE conjE)
  assume len: length (list-of ltrv2') < length (list-of ltrv2)
  show ?thesis
  apply(rule less(1)[OF - ltrv2'])
  subgoal by fact subgoal by fact
  subgoal using  $\omega4$  unfolding  $\omega4$ -def by simp
  subgoal by fact
  subgoal using r(2)  $\omega\omega$  by (metis Opt.reach.Step fst-conv snd-conv)
  subgoal using  $\omega4$  unfolding  $\omega4$ -def
  by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(3))

```

```

      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
      subgoal by fact subgoal by fact subgoal by fact subgoal by fact
      subgoal by fact subgoal by fact .
next
assume len: length (list-of ltrv2') = length (list-of ltrv2) w2' < w2
show ?thesis
apply(rule less(2)[OF - - ltrv2'])
      subgoal by fact subgoal using len by simp subgoal by fact

      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
      subgoal by fact
      subgoal by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(2) snd-conv)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2 r(4))
      subgoal by fact subgoal by fact subgoal by fact subgoal by fact
      subgoal by fact subgoal by fact .
qed
show ?thesis unfolding lltrv2 using 0
by (simp add: lfin' list-of-lappend)
next
case False note current = current False
obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
XX: XX s2 ltr2 tr2 s2' s2'' ltr2' and
 $\chi_4'$ :  $\chi_4'$   $\Delta$  w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and ltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF unW  $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]
  unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2' = lltrv2 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
  have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding ltrv2 ltrv2' ..

  have trv2ne: trv2  $\neq$  [] using  $\chi_4'$  unfolding  $\chi_4'$ -def by auto
  have lfin': lfinite ltrv2'
  using lfin trv2ne unfolding lltrv2 by simp
  have len: length (list-of ltrv2') < length (list-of ltrv2)
  using trv2ne lfin lfin' by (simp add: list-of-lappend lltrv2)

  have 0: list-of ltrv2'  $\neq$  []  $\wedge$  finalV (last (list-of ltrv2'))
  apply(rule less(1)[OF - ltrv2'])

```

```

    subgoal by fact subgoal by fact
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
    subgoal by fact
    subgoal using  $r(2)$   $\chi\chi$  unfolding  $\chi\chi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2  $r(3)$ )
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
 $r(4)$ )
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast .
    show ?thesis unfolding lltrv2 using 0
      by (simp add: lfn' list-of-lappend)
  qed
qed
qed
qed
qed
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

```

lemma *lS-lltrv1-ltr1*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$

and *r*: *reachO* *s1 reachO s2 reachV sv1 reachV sv2*

and *ltr1*: *Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1*

and *ltr2*: *Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2*

shows *Van.lS (lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) = Opt.lS ltr1*

proof–

have *cltrv1*: *Van.lcompletedFrom sv1 (lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))*

using *lcompletedFrom-lltrv1[OF assms]* .

{fix *trn nL nR ltrv1 ltr1*

assume $\exists w1 w2 s1 s2 ltr2 statA sv1 sv2 statO$.

$nL = w1 \wedge nR = w2 \wedge$

$ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$

$\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$

$reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$

$Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge lnever isIntO ltr1 \wedge$

$Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge lnever isIntO ltr2$

hence *TwoFuncPred.sameFM1 isSecV isSecO getSecV getSecO trn nL nR ltrv1*

ltr1

proof(*coinduct rule: TwoFuncPred.sameFM1.coinduct*[of $\lambda trn\ nL\ nR\ ltrv1\ ltr1.$

$\exists w1\ w2\ s1\ s2\ ltr2\ statA\ sv1\ sv2\ statO.$
 $nL = w1 \wedge nR = w2 \wedge$
 $ltrv1 = lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$
 $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge lnever\ isIntO\ ltr1 \wedge$

$Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge lnever\ isIntO\ ltr2,$
where $pred = isSecV$ **and** $pred' = isSecO$ **and** $func = getSecV$ **and** $func' = getSecO$])

case ($2\ trn\ nL\ nR\ ltrv1\ ltr1$)

then obtain $w1\ w2\ sv1\ s1\ s2\ ltr2\ statA\ sv2\ statO$

where $nL: nL = w1$ **and** $nR: nR = w2$

and $ltrv1: ltrv1 = lltrv1\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$

and $\Delta: \Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$

and $ltr1: Opt.lvalidFromS\ s1\ ltr1\ Opt.lcompletedFrom\ s1\ ltr1\ lnever\ isIntO\ ltr1$

and $ltr2: Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2\ lnever\ isIntO\ ltr2$

by *auto*

have $isi3: \neg isIntO\ s1$ **using** $ltr1$

by (*metis* $Opt.lcompletedFrom-def\ Opt.lvalidFromS-def\ lfinite-LNil\ llist.exhaust-sel\ llist.pred-inject(2)$)

have $isi4: \neg isIntO\ s2$ **using** $ltr2$

by (*metis* $Opt.lcompletedFrom-def\ Opt.lvalidFromS-def\ lfinite-LNil\ llist.exhaust-sel\ llist.pred-inject(2)$)

show $?case\ \mathbf{proof}(cases\ ltr1 = [] \wedge ltr2 = [])$

case $True$ **note** $ltr14 = True$

hence $ltrv1: ltrv1 = []$ **unfolding** $ltrv1$ **by** *simp*

show $?thesis$ **using** $ltr14$ **unfolding** $ltrv1$ **apply-apply**(*rule* $TwoFuncPred.sameFM1-selectLNil$)

by *auto*

next

case $False$ **hence** $ltr14: ltr1 \neq [] \vee ltr2 \neq []$ **by** *auto*

show $?thesis\ \mathbf{proof}(cases\ llength\ ltr1 \leq Suc\ 0 \vee llength\ ltr2 \leq Suc\ 0)$

case $True$ **note** $ltr14 = ltr14\ True$

hence $ltrv1: ltrv1 = [[sv1]]$ **unfolding** $ltrv1$ **by** *simp*

have $llength\ ltr1 = Suc\ 0 \vee llength\ ltr2 = Suc\ 0$

by (*metis* $Opt.lcompletedFrom-def\ Suc-ile-eq\ True$

$lfinite-LNil\ llength-LNil\ llist-eq-cong\ ltr1(2)$

$ltr2(2)\ nle-le\ order-le-imp-less-or-eq\ zero-enat-def\ zero-order(3)$)

hence $finalO\ s1 \vee finalO\ s2$

using $Opt.lcompletedFrom-singl\ ltr1(1)\ ltr1(2)\ ltr2(1)\ ltr2(2)$ **by** *blast*

hence $fs1: finalO\ s1$

using $\Delta\ r(1)\ r(2)\ r(3)\ r(4)\ unw\ unwindCond-def$ **by** *auto*

hence $ltr1: ltr1 = [[s1]]$

by (*metis* $Opt.final-def\ Opt.lcompletedFrom-def$

```

      Opt.lvalidFromS-Cons-iff lfinite-code(1) llist.exhaust ltr1(1) ltr1(2))
    have fsv1: finalV sv1
    using  $\Delta$  fs1 r(1) r(2) r(3) r(4) unW unwindCond-final by blast
    have isv13:  $\neg$  isSecV sv1  $\wedge$   $\neg$  isSecO s1
    using fsv1 fs1 Opt.final-not-isSec Van.final-not-isSec by blast
  show ?thesis unfolding ltrv1 ltr1 apply(rule TwoFuncPred.sameFM1-selectSingl)

  using isv13 by auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
by auto
show ?thesis proof(cases trn)
  case L note trn = L[simp] note current = current L
  show ?thesis
  proof(cases lnever isSecO ltr1)
    case True note current = current True
    obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
       $\omega\omega$ : ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
      lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
       $\omega\omega$ :  $\omega\omega$   $\Delta$  w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
  statOO
    and trn': trn' = (if trv1 = [] then L else R)
    and lltrv1: ltrv1 =
      lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
  sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unW  $\Delta$  r ltr1 isi4 current]
    unfolding ltrv1 by blast
    define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (trn', w1', w2', s1', ltr1',
  s2, ltr2, statA, sv1', sv2', statOO)
    have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
    unfolding lltrv1 ltrv1' ..
    have nis1:  $\neg$  isSecO s1 using True  $\omega\omega$ (1) by force
    show ?thesis
    proof(cases trv1 = [])
      case True note trv1 = True
      hence w1' < w1 using  $\omega\omega$  unfolding  $\omega\omega$ -def by auto
      have [simp]: trn' = trn by (simp add: trv1 trn')
      show ?thesis
      apply(rule TwoFuncPred.sameFM1-selectDelayL)
      apply(rule exI[of - w1']) apply(rule exI[of - w1])
      apply(rule exI[of - trv1]) apply(rule exI[of - [s1]])
      apply(rule exI[of - w2'])
      apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1'])
      apply(rule exI[of - w2])
      apply(intro conjI)
      subgoal by fact
      subgoal unfolding nL .. subgoal unfolding nR ..
      subgoal unfolding ltrv1 trv1 by simp
      subgoal unfolding  $\omega\omega$ (1) by simp

```

```

subgoal by fact subgoal unfolding trv1 using  $\omega 3$ -def nis1 by
simp
subgoal apply(rule disjI1)
  apply(rule exI[of - w1]) apply(rule exI[of - w2])
  apply(rule exI[of - s1]) apply(rule exI[of - s2])
  apply(rule exI[of - ltr2]) apply(rule exI[of - statA])
  apply(rule exI[of - sv1]) apply(rule exI[of - sv2])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3))
snoc-eq-iff-butlast)
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4))
snoc-eq-iff-butlast)
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\omega\omega$  using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
next
case False note trv1 = False
show ?thesis
apply(rule TwoFuncPred.sameFM1-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - [s1]])
apply(rule exI[of - trn]) apply(rule exI[of - w1])
apply(rule exI[of - w2])
apply(rule exI[of - ltrv1]) apply(rule exI[of - ltr1])
apply(rule exI[of - trn])
apply(rule exI[of - w1])
apply(rule exI[of - w2])
apply(intro conjI)
  subgoal ..
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv1 .
  subgoal unfolding  $\omega\omega$ (1) by simp
  subgoal by fact
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
  subgoal using ltr1(3)  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Opt.S.map-filter Opt.S.simps(4) Van.S.map-filter
Van.S.eq-Nil-iff(2) append-Nil

```

```

butlast-snoc filter.simps(2) nis1
subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1']) apply(rule exI[of - s2'])
  apply(rule exI[of - ltr2']) apply(rule exI[of - statA])
  apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' ..
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\omega\omega$  using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .

qed
next
case False note current = current False
obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
 $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
 $\chi\mathcal{P}'$ :  $\chi\mathcal{P}' \Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and lltrv1: ltrv1 =
lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

using lltrv1-lltrv2-not-lnever-L[OF unW  $\Delta$  r ltr1 isi4 current]
unfolding ltrv1 by blast
define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
unfolding lltrv1 ltrv1' ..

show ?thesis apply(rule TwoFuncPred.sameFM1-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - tr1 ## s1'])
apply(rule exI[of - R])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - s1'' $ ltr1'])

```

```

apply(rule exI[of - trn])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal .. subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv1 .
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by simp
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by simp
  subgoal by simp
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
  by (simp add: Opt.S.map-filter Van.S.map-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1'])
  apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule exI[of
- sv2'])
  apply(rule exI[of - statOO])
  apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrv1' ..
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by simp
    subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

  append-is-Nil-conv last-snoc not-Cons-self2 r(1))
  subgoal by fact
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
qed
next
case R note trn = R[simp] note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
case True note current = current True
obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
 $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
 $\omega\omega$ :  $\omega\omega$   $\Delta$  w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'

```

```

statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv1: ltrv1 =
    lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF unW Δ r ltr2(1,2) isi3 ltr2(3) current]
  unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..
  have nev1: never isSecV trv1 using ω4 unfolding ω4-def by auto
  show ?thesis
  proof(cases trv2 = [])
    case True note trv2 = True
    have [simp]: trn' = trn using R trv2 trn' by auto
    have w2' < w2 using ω4 trv2 unfolding ω4-def by auto
    show ?thesis
    apply(rule TwoFuncPred.sameFM1-selectDelayR)
    apply(rule exI[of - w2']) apply(rule exI[of - nR])
    apply(rule exI[of - trv1]) apply(rule exI[of - []])
    apply(rule exI[of - w1'])
    apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1])
    apply(rule exI[of - nL])
    apply(intro conjI)
    subgoal by simp subgoal .. subgoal ..
    subgoal by fact subgoal by simp
    subgoal unfolding nR by fact
    subgoal using nev1 by (simp add: never-Nil-filter)
    subgoal apply(rule disjI1)
      apply(rule exI[of - w1']) apply(rule exI[of - w2'])
      apply(rule exI[of - s1]) apply(rule exI[of - s2'])
      apply(rule exI[of - ltr2']) apply(rule exI[of - statA])
      apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
      apply(rule exI[of - statOO])
      apply(intro conjI)
      subgoal .. subgoal ..
      subgoal unfolding ltrv1' by simp
      subgoal using ω4 unfolding ω4-def by simp
      subgoal by fact
      subgoal using ω4 unfolding ω4-def
    by (metis Opt.reach.Step ωω(2) fst-conv r(2) snd-conv)
    subgoal using ω4 unfolding ω4-def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using ω4 unfolding ω4-def
    by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
    subgoal by fact subgoal by fact subgoal by fact
    subgoal using ωω by auto
    subgoal using ωω by auto

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```

      subgoal using  $\omega\omega$  by auto . .
next
case False note trv2 = False
have [simp]: trn' = L using R trv2 trn' by auto
show ?thesis
apply(rule TwoFuncPred.sameFM1-selectRL)
apply(rule exI[of - trv1]) apply(rule exI[of - []])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal by fact
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal unfolding ltrv1 ..
  subgoal unfolding  $\omega\omega(1)$  by simp
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by (simp add: never-Nil-filter)
subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1]) apply(rule exI[of - s2'])
  apply(rule exI[of - ltr2']) apply(rule exI[of - statA])
  apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv1' by simp
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
subgoal by fact subgoal by fact subgoal by fact
subgoal using  $\omega\omega$  by auto
subgoal using  $\omega\omega$  by auto
subgoal using  $\omega\omega$  by auto . .
qed
next
case False note current = current False
obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where
 $\chi\chi$ :  $\chi\chi$  s2 ltr2 tr2 s2' s2'' ltr2' and
 $\chi_4'$ :  $\chi_4' \Delta$  w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (L, w1', w2', s1, ltr1, s2'', s2'' $ ltr2',
statA, sv1'', sv2'', statOO))
  using lltrv1-lltrv2-not-lnever-R[OF un $\omega$   $\Delta$  r ltr2(1,2) isi3 ltr2(3)
current]

```

```

unfolding ltrv1 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (L, w1', w2', s1, ltr1, s2'',
s2'' $ ltr2', statA, sv1'', sv2'', statOO)
  have ltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding ltrv1 ltrv1' ..

show ?thesis
apply(rule TwoFuncPred.sameFM1-selectRL)
apply(rule exI[of - trv1]) apply(rule exI[of - []])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv1']) apply(rule exI[of - ltr1'])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(intro conjI)
  subgoal by fact
    subgoal unfolding nL .. subgoal unfolding nR ..
    subgoal unfolding ltrv1 .. subgoal by simp
subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by (simp add: never-Nil-filter)
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - s2''])
    apply(rule exI[of - s2'' $ ltr2']) apply(rule exI[of - statA'])
    apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
    apply(rule exI[of - statOO'])
    apply(intro conjI)
      subgoal .. subgoal ..
      subgoal unfolding ltrv1' by simp
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
      subgoal by fact
      subgoal using r(2)  $\chi\chi$  unfolding  $\chi\chi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
      subgoal using r(3)  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal using r(4)  $\chi_4'$  unfolding  $\chi_4'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast . .

    qed
  qed
qed
qed
}
thus ?thesis unfolding Van.lS[OF cltrv1] Opt.lS[OF ltr1(2)]
apply– apply(rule TwoFuncPred.sameFM1-lmap-lfilter)
using assms by blast

```


qed

lemma *lS-lltrv2-ltr2*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever isIntO* *ltr1*

and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever isIntO* *ltr2*

shows *Van.lS* (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)) = *Opt.lS* *ltr2*

proof –

have *cltrv2*: *Van.lcompletedFrom* *sv2* (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

using *lcompletedFrom-lltrv2* [*OF* *assms*]

{fix *trn* *nL* *nR* *ltrv2* *ltr2*

assume $\exists w1\ w2\ s1\ s2\ ltr1\ statA\ sv1\ sv2\ statO.$

$nL = w1 \wedge nR = w2 \wedge$

$ltrv2 = lltrv2\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$

$reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$

$Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge lnever\ isIntO\ ltr1 \wedge$

$Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge lnever\ isIntO\ ltr2$

hence *TwoFuncPred.sameFM2* *isSecV* *isSecO* *getSecV* *getSecO* *trn* *nL* *nR* *ltrv2* *ltr2*

proof (*coinduct* rule: *TwoFuncPred.sameFM2.coinduct* [of $\lambda trn\ nL\ nR\ ltrv2\ ltr2.$

$\exists w1\ w2\ s1\ s2\ ltr1\ statA\ sv1\ sv2\ statO.$

$nL = w1 \wedge nR = w2 \wedge$

$ltrv2 = lltrv2\ (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$

$reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$

$Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge lnever\ isIntO\ ltr1 \wedge$

$Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge lnever\ isIntO\ ltr2,$

where *pred* = *isSecV* **and** *pred'* = *isSecO* **and** *func* = *getSecV* **and** *func'* = *getSecO*])

case (*2* *trn* *nL* *nR* *ltrv2* *ltr2*)

then obtain *w1* *w2* *sv1* *s1* *s2* *ltr1* *statA* *sv2* *statO*

where *nL*: *nL* = *w1* **and** *nR*: *nR* = *w2*

and *ltrv2*: *ltrv2* = *lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)

and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever isIntO* *ltr1*

and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever isIntO* *ltr2*

by *auto*

have *isi3*: $\neg isIntO\ s1$ **using** *ltr1*

by (*metis* *Opt.lcompletedFrom-def* *Opt.lvalidFromS-def* *lfinite-LNil* *l**list.exhaust-sel* *l**list.pred-inject* (*2*))

have *isi4*: $\neg isIntO\ s2$ **using** *ltr2*

by (metis *Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel llist.pred-inject(2)*)

```

show ?case proof(cases ltr1 = [] ^ ltr2 = [])
  case True note ltr14 = True
  hence ltrv2: ltrv2 = [] unfolding ltrv2 by simp
show ?thesis using ltr14 unfolding ltrv2 apply-apply(rule TwoFuncPred.sameFM2-selectLNil)
by auto
next
case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
  hence ltrv2: ltrv2 = [[sv2]] unfolding ltrv2 by simp
  have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
  by (metis Opt.lcompletedFrom-def Suc-ile-eq True lfinite-LNil llength-LNil llist-eq-cong ltr1(2) ltr2(2) nle-le order-le-imp-less-or-eq zero-enat-def zero-order(3))
  hence finalO s1 ∨ finalO s2
  using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
  hence fs2: finalO s2
  using Δ r(1) r(2) r(3) r(4) unW unwindCond-def by auto
  hence ltr2: ltr2 = [[s2]]
  by (metis Opt.final-def Opt.lcompletedFrom-def Opt.lvalidFromS-Cons-iff lfinite-code(1) llist.exhaust ltr2(1) ltr2(2))
  have fsv2: finalV sv2
  using Δ fs2 r(1) r(2) r(3) r(4) unW unwindCond-final by blast
  have isv24: ¬ isSecV sv2 ∧ ¬ isSecO s2
  using fsv2 fs2 Opt.final-not-isSec Van.final-not-isSec by blast
show ?thesis unfolding ltrv2 ltr2 apply(rule TwoFuncPred.sameFM2-selectSingl)

  using isv24 by auto
next
case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0
by auto
show ?thesis proof(cases trn)
  case L note trn = L[simp] note current = current L
  show ?thesis
  proof(cases lnever isSecO ltr1)
    case True note current = current True
    obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
      ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1' lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
      ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
      statOO
    and trn': trn' = (if trv1 = [] then L else R)
    and lltrv2: ltrv2 =
      lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
      sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unW Δ r ltr1 isi4 current]
  
```

```

unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
  have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding lltrv2 ltrv2' ..
  have nev2: never isSecV trv2 using  $\omega 3$  unfolding  $\omega 3$ -def by auto
  show ?thesis
proof(cases trv1 = [])
  case True note trv1 = True
  have [simp]: trn' = trn using L trv1 trn' by auto
  have w1' < w1 using  $\omega 3$  trv1 unfolding  $\omega 3$ -def by auto
  show ?thesis
  apply(rule TwoFuncPred.sameFM2-selectDelayL)
  apply(rule exI[of - w1 ^]) apply(rule exI[of - nL])
  apply(rule exI[of - trv2]) apply(rule exI[of - []])
  apply(rule exI[of - w2 ^])
  apply(rule exI[of - ltrv2 ^]) apply(rule exI[of - ltr2])
  apply(rule exI[of - nR])
  apply(intro conjI)
  subgoal by simp subgoal .. subgoal ..
  subgoal by fact subgoal by simp
  subgoal unfolding nL by fact
  subgoal using nev2 by (simp add: never-Nil-filter)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1 ^]) apply(rule exI[of - w2 ^])
  apply(rule exI[of - s1 ^]) apply(rule exI[of - s2])
  apply(rule exI[of - ltr1 ^]) apply(rule exI[of - statA])
  apply(rule exI[of - sv1 ^]) apply(rule exI[of - sv2 ^])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv2' by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def by simp
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
  subgoal using  $\omega 3$  unfolding  $\omega 3$ -def
  by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal by fact subgoal by fact subgoal by fact ..
next
  case False note trv1 = False
  have [simp]: trn' = R using L trv1 trn' by auto
  show ?thesis
  apply(rule TwoFuncPred.sameFM2-selectLR)

```

```

apply(rule exI[of - trv2]) apply(rule exI[of - []])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal by fact
    subgoal unfolding nL .. subgoal unfolding nR ..
    subgoal unfolding ltrv2 ..
    subgoal unfolding  $\omega\omega(1)$  by simp
subgoal using  $\omega\beta$  unfolding  $\omega\beta$ -def by (simp add: never-Nil-filter)
subgoal apply(rule disjI1)
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1']) apply(rule exI[of - s2])
apply(rule exI[of - ltr1']) apply(rule exI[of - statA])
apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal .. subgoal ..
    subgoal unfolding ltrv2' by simp
    subgoal using  $\omega\beta$  unfolding  $\omega\beta$ -def by simp
    subgoal using  $\omega\beta$  unfolding  $\omega\beta$ -def
    by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(1) snd-conv)
    subgoal by fact
    subgoal using  $\omega\beta$  unfolding  $\omega\beta$ -def
    by (metis Van.reach-validFromS-reach r(3) snoc-eq-iff-butlast)
    subgoal using  $\omega\beta$  unfolding  $\omega\beta$ -def
    by (metis Van.reach-validFromS-reach r(4) snoc-eq-iff-butlast)
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal using  $\omega\omega$  by auto
    subgoal by fact subgoal by fact subgoal by fact ..
qed
next
case False note current = current False
obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
 $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
 $\chi\beta$ :  $\chi\beta$   $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
and lltrv2: ltrv2 =
lappend (llist-of trv2) (lltrv2 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

using lltrv1-lltrv2-not-lnever-L[OF un $\omega$   $\Delta$  r ltr1 isi4 current]
unfolding ltrv2 by blast
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
unfolding lltrv2 ltrv2' ..

show ?thesis

```

```

apply(rule TwoFuncPred.sameFM2-selectLR)
apply(rule exI[of - trv2]) apply(rule exI[of - []])
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2'])
apply(rule exI[of - w1]) apply(rule exI[of - w2])
apply(intro conjI)
  subgoal by fact
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal unfolding ltrv2 .. subgoal by simp
subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by (simp add: never-Nil-filter)
subgoal apply(rule disjI1)
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1'']) apply(rule exI[of - s2'])
apply(rule exI[of - s1'' $ ltr1']) apply(rule exI[of - statA])
apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv2' by simp
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'-def}$  by simp
  subgoal using r(1)  $\chi\chi$  unfolding  $\chi\chi-def$ 
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal by fact
  subgoal using r(3)  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using r(4)  $\chi^{3'}$  unfolding  $\chi^{3'-def}$ 
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$  by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi-def$ 
    using l1ist-all-lappend-l1ist-of ltr1(3) by blast
  subgoal by fact subgoal by fact subgoal by fact . .
qed
next
case R note trn = R[simp] note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where
     $\omega\omega$ : ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
    lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
     $\omega4$ :  $\omega4 \Delta w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'$ 
statOO
  and trn': trn' = (if trv2 = [] then R else L)
  and ltrv2: ltrv2 =
    lappend (l1ist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
  using lltrv1-lltrv2-lnever-R[OF unW \Delta r ltr2(1,2) isi3 ltr2(3) current]

```

```

unfolding ltrv2 by blast
define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
have ltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
unfolding ltrv2 ltrv2' ..
have nis2:  $\neg$  isSecO s2 using True  $\omega\omega(1)$  by force

show ?thesis
proof(cases trv2 = [])
case True note trv2 = True
hence w2' < w2 using  $\omega_4$  unfolding  $\omega_4$ -def by auto
have [simp]: trn' = trn by (simp add: trv2 trn')
show ?thesis
apply(rule TwoFuncPred.sameFM2-selectDelayR)
apply(rule exI[of - w2']) apply(rule exI[of - w2])
apply(rule exI[of - trv2]) apply(rule exI[of - [s2]])
apply(rule exI[of - w1'])
apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2'])
apply(rule exI[of - w1])
apply(intro conjI)
subgoal by fact
subgoal unfolding nL .. subgoal unfolding nR ..
subgoal unfolding ltrv2 trv2 by simp
subgoal unfolding  $\omega\omega(1)$  by simp
subgoal by fact subgoal unfolding trv2 using  $\omega_4$ -def nis2 by
simp
subgoal apply(rule disjI1)
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1]) apply(rule exI[of - s2'])
apply(rule exI[of - ltr1]) apply(rule exI[of - statA])
apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
subgoal .. subgoal ..
subgoal unfolding ltrv2' by simp
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
subgoal by fact
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Simple-Transition-System.reach-validFromS-reach r(3))
snoc-eq-iff-butlast)
subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
by (metis Simple-Transition-System.reach-validFromS-reach r(4))
snoc-eq-iff-butlast)
subgoal by fact subgoal by fact subgoal by fact
subgoal using  $\omega\omega$  by auto
subgoal using  $\omega\omega$  by auto
subgoal using  $\omega\omega$ 

```

```

        using llist-all-lappend-llist-of ltr1(3) by blast . .
next
  case False note trv2 = False
  show ?thesis
  apply(rule TwoFuncPred.sameFM2-selectlappend)
  apply(rule exI[of - trv2]) apply(rule exI[of - [s2]])
  apply(rule exI[of - trn']) apply(rule exI[of - w1'])
  apply(rule exI[of - w2'])
  apply(rule exI[of - ltrv2']) apply(rule exI[of - ltr2'])
  apply(rule exI[of - trn])
  apply(rule exI[of - w1])
  apply(rule exI[of - w2])
  apply(intro conjI)
  subgoal ..
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv2 .
  subgoal unfolding  $\omega\omega(1)$  by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal using ltr1(3)  $\omega_4$  unfolding  $\omega_4$ -def
  by (simp add: never-Nil-filter nis2)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1]) apply(rule exI[of - s2'])
  apply(rule exI[of - ltr1]) apply(rule exI[of - statA])
  apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
  apply(rule exI[of - statOO])
  apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltrv2' ..
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def by simp
  subgoal by fact
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal by fact subgoal by fact subgoal by fact
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast . .
qed
next
  case False note current = current False
  obtain w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO where

```

```

     $\chi\chi$ :  $\chi\chi$   $s2$   $ltr2$   $tr2$   $s2'$   $s2''$   $ltr2'$  and
     $\chi4'$ :  $\chi4'$   $\Delta$   $w1$   $w2$   $w1'$   $w2'$   $s1$   $s2$   $tr2$   $s2'$   $s2''$   $statA$   $sv1$   $trv1$   $sv1''$   $sv2$ 
     $trv2$   $sv2''$   $statOO$ 
    and  $ltrv2$ :  $ltrv2$  =
     $lappend$  ( $l$ list-of  $trv2$ ) ( $lltrv2$  ( $L$ ,  $w1'$ ,  $w2'$ ,  $s1$ ,  $ltr1$ ,  $s2''$ ,  $s2''$   $\$$   $ltr2'$ ,
     $statA$ ,  $sv1''$ ,  $sv2''$ ,  $statOO$ ))
    using  $lltrv1$ - $lltrv2$ -not-lnever-R[ $OF$   $unw$   $\Delta$   $r$   $ltr2(1,2)$   $isi3$   $ltr2(3)$ 
     $current$ ]
    unfolding  $ltrv2$  by  $blast$ 
    define  $ltrv2'$  where  $ltrv2'$ :  $ltrv2' \equiv lltrv2$  ( $L$ ,  $w1'$ ,  $w2'$ ,  $s1$ ,  $ltr1$ ,  $s2''$ ,
     $s2''$   $\$$   $ltr2'$ ,  $statA$ ,  $sv1''$ ,  $sv2''$ ,  $statOO$ )
    have  $ltrv2$ :  $ltrv2 = lappend$  ( $l$ list-of  $trv2$ )  $ltrv2'$ 
    unfolding  $ltrv2$   $ltrv2'$  ..
    show  $?thesis$ 
    apply( $rule$   $TwoFuncPred.sameFM2-selectlappend$ )
    apply( $rule$   $exI$ [of -  $trv2$ ]) apply( $rule$   $exI$ [of -  $tr2$   $##$   $s2'$ ])
    apply( $rule$   $exI$ [of -  $L$ ])
    apply( $rule$   $exI$ [of -  $w1$ ]) apply( $rule$   $exI$ [of -  $w2$ ])
    apply( $rule$   $exI$ [of -  $ltrv2$ ]) apply( $rule$   $exI$ [of -  $s2''$   $\$$   $ltr2'$ ])
    apply( $rule$   $exI$ [of -  $trn$ ])
    apply( $rule$   $exI$ [of -  $w1$ ]) apply( $rule$   $exI$ [of -  $w2$ ])
    apply( $intro$   $conjI$ )
    subgoal .. subgoal unfolding  $nL$  .. subgoal unfolding  $nR$  ..
    subgoal using  $ltrv2$  .
    subgoal using  $\chi\chi$  unfolding  $\chi\chi$ - $def$  by  $simp$ 
    subgoal using  $\chi4'$  unfolding  $\chi4'$ - $def$  by  $simp$ 
    subgoal by  $simp$ 
    subgoal using  $\chi4'$  unfolding  $\chi4'$ - $def$ 
    by ( $simp$   $add$ :  $Opt.S.map-filter$   $Van.S.map-filter$ )
    subgoal apply( $rule$   $disjI1$ )
    apply( $rule$   $exI$ [of -  $w1$ ]) apply( $rule$   $exI$ [of -  $w2$ ])
    apply( $rule$   $exI$ [of -  $s1$ ])
    apply( $rule$   $exI$ [of -  $s2''$ ]) apply( $rule$   $exI$ [of -  $ltr1$ ])
    apply( $rule$   $exI$ [of -  $statA$ ]) apply( $rule$   $exI$ [of -  $sv1''$ ]) apply( $rule$   $exI$ [of
    -  $sv2''$ ])

    apply( $rule$   $exI$ [of -  $statOO$ ])
    apply( $intro$   $conjI$ )
    subgoal .. subgoal ..
    subgoal unfolding  $ltrv2'$  ..
    subgoal using  $\chi4'$  unfolding  $\chi4'$ - $def$  by  $simp$ 
    subgoal by  $fact$ 
    subgoal using  $\chi4'$  unfolding  $\chi4'$ - $def$ 
by ( $metis$   $Simple-Transition-System.reach-validFromS-reach$   $\chi\chi$   $\chi\chi$ - $def$ 

     $append-is-Nil-conv$   $last-snoc$   $not-Cons-self2$   $r(2)$ )
    subgoal using  $\chi4'$  unfolding  $\chi4'$ - $def$ 
    by ( $metis$   $Van.reach-validFromS-reach$   $r(3)$   $snoc-eq-iff-butlast$ )
    subgoal using  $\chi4'$  unfolding  $\chi4'$ - $def$ 
    by ( $metis$   $Van.reach-validFromS-reach$   $r(4)$   $snoc-eq-iff-butlast$ )

```



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      subgoal by fact subgoal by fact subgoal by fact
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
      subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
      using llist-all-lappend-llist-of ltr2(3) by blast . .
    qed
  qed
  qed
  qed
  }
  thus ?thesis unfolding Van.lS[OF cltrv2] Opt.lS[OF ltr2(2)]
  apply- apply(rule TwoFuncPred.sameFM2-lmap-lfilter)
  using assms by blast
qed

```

lemma *lA-lltrv1-lltrv2*:

assumes *unw*: *unwindCond* Δ

and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$

and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*

and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1* *lnever isIntO* *ltr1*

and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2* *lnever isIntO* *ltr2*

shows *Van.lA* (*lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*)) =

Van.lA (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

proof–

have *cltrv1*: *Van.lcompletedFrom* *sv1* (*lltrv1* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

using *lcompletedFrom-lltrv1*[*OF* *assms*] .

have *cltrv2*: *Van.lcompletedFrom* *sv2* (*lltrv2* (*trn*, *w1*, *w2*, *s1*, *ltr1*, *s2*, *ltr2*, *statA*, *sv1*, *sv2*, *statO*))

using *lcompletedFrom-lltrv2*[*OF* *assms*] .

{fix *nL* *nR* *ltrv1* *ltrv2*

assume $\exists trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO$.

$nL = w1 \wedge nR = w2 \wedge$

$ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$ltrv2 = lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$

$reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$

$Opt.lvalidFromS s1 ltr1 \wedge Opt.lcompletedFrom s1 ltr1 \wedge lnever isIntO ltr1 \wedge$

$Opt.lvalidFromS s2 ltr2 \wedge Opt.lcompletedFrom s2 ltr2 \wedge lnever isIntO ltr2$

hence *TwoFuncPred.sameFM isIntV isIntV getActV getActV* *nL* *nR* *ltrv1* *ltrv2*

proof(*coinduct* rule: *TwoFuncPred.sameFM.coinduct*[*of* $\lambda nL nR ltrv1 ltrv2$].

$\exists trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO$.

$nL = w1 \wedge nR = w2 \wedge$

$ltrv1 = lltrv1 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$ltrv2 = lltrv2 (trn, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) \wedge$

$\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$

$reachO s1 \wedge reachO s2 \wedge reachV sv1 \wedge reachV sv2 \wedge$

```

    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧ lnever isIntO ltr1 ∧
    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧ lnever isIntO ltr2])
  case (2 nL nR ltrv1 ltrv2)
  then obtain trn w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
  where nL: nL = w1 and nR: nR = w2
  and ltrv1: ltrv1 = lltrv1 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and ltrv2: ltrv2 = lltrv2 (trn,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1 lnever isIntO ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2 lnever isIntO ltr2
  by auto
  have isi3: ¬ isIntO s1 using ltr1
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
  llist.pred-inject(2))
  have isi4: ¬ isIntO s2 using ltr2
  by (metis Opt.lcompletedFrom-def Opt.lvalidFromS-def lfinite-LNil llist.exhaust-sel
  llist.pred-inject(2))

  show ?case proof(cases ltr1 = [] ∧ ltr2 = [])
  case True note ltr14 = True
  hence ltrv1: ltrv1 = [] unfolding ltrv1 by simp
  show ?thesis using ltr14 unfolding ltrv1 ltrv2 apply-apply(rule Two-
  FuncPred.sameFM-selectLNil) by auto
  next
  case False hence ltr14: ltr1 ≠ [] ∨ ltr2 ≠ [] by auto
  show ?thesis proof(cases llength ltr1 ≤ Suc 0 ∨ llength ltr2 ≤ Suc 0)
  case True note ltr14 = ltr14 True
  hence ltrv1: ltrv1 = [[sv1]] and ltrv2: ltrv2 = [[sv2]] unfolding ltrv1 ltrv2
  by auto
  have llength ltr1 = Suc 0 ∨ llength ltr2 = Suc 0
  by (metis Opt.lcompletedFrom-def Suc-ile-eq True
  lfinite-LNil llength-LNil llist-eq-cong ltr1(2)
  ltr2(2) nle-le order-le-imp-less-or-eq zero-enat-def zero-order(3))
  hence finalO s1 ∨ finalO s2
  using Opt.lcompletedFrom-singl ltr1(1) ltr1(2) ltr2(1) ltr2(2) by blast
  hence fs1: finalO s1 ∧ finalO s2
  using Δ r(1) r(2) r(3) r(4) unw unwindCond-def by auto

  have fsv12: finalV sv1 ∧ finalV sv2
  using Δ fs1 r(1) r(2) r(3) r(4) unw unwindCond-final by blast
  have isv12: ¬ isIntV sv1 ∧ ¬ isIntV sv2
  using fsv12 Van.final-not-isInt by blast
  show ?thesis unfolding ltrv1 ltrv2 apply(rule TwoFuncPred.sameFM-selectSingl)

  using isv12 by auto
  next
  case False hence current: llength ltr1 > Suc 0 llength ltr2 > Suc 0

```

```

by auto
show ?thesis proof(cases trn)
  case L note current = current L
  show ?thesis
  proof(cases lnever isSecO ltr1)
    case True note current = current True
    obtain trn' w1' w2' s1' ltr1' trv1 sv1' trv2 sv2' statOO where
      ωω: ltr1 = s1 $ ltr1' validTransO (s1, s1') Opt.lvalidFromS s1' ltr1'
      lcompletedFromO s1' ltr1' lnever isIntO ltr1' and
      ω3: ω3 Δ w1 w2 w1' w2' s1 s1' s2 statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
    and trn': trn' = (if trv1 = [] then L else R)
    and lltrv1: ltrv1 =
      lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    and lltrv2: ltrv2 =
      lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1', ltr1', s2, ltr2, statA,
sv1', sv2', statOO))
    using lltrv1-lltrv2-lnever-L[OF unω Δ r ltr1 isi4 current]
    unfolding ltrv1 ltrv2 by blast
    define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
    have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
    unfolding lltrv1 ltrv1' ..
    define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (trn', w1', w2', s1', ltr1',
s2, ltr2, statA, sv1', sv2', statOO)
    have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
    unfolding lltrv2 ltrv2' ..

    show ?thesis
    apply(rule TwoFuncPred.sameFM-selectlappend)
    apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of -
w1])
    apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of -
w2])
    apply(rule exI[of - ltrv1']) apply(rule exI[of - ltrv2'])
    apply(intro conjI)
    subgoal unfolding nL .. subgoal unfolding nR ..
    subgoal using lltrv1 .
    subgoal using lltrv2 .
    subgoal using ω3 unfolding ω3-def by simp
    subgoal using ω3 unfolding ω3-def by simp
    subgoal using ω3 unfolding ω3-def by (simp add: Van.A.map-filter)

    subgoal apply(rule disjI1)
    apply(rule exI[of - trn']) apply(rule exI[of - w1']) apply(rule exI[of
- w2'])
    apply(rule exI[of - s1']) apply(rule exI[of - ltr1'])
    apply(rule exI[of - s2]) apply(rule exI[of - ltr2])

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apply(rule exI[of - statA])
apply(rule exI[of - sv1']) apply(rule exI[of - sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal .. subgoal ..
  subgoal unfolding ltr1' ..
  subgoal unfolding ltr2' ..
  subgoal using  $\omega_3$  unfolding  $\omega_3$ -def by simp
  subgoal using  $\omega_3$  unfolding  $\omega_3$ -def
  by (metis Opt.reach.Step  $\omega\omega$ (2) fst-conv r(1) snd-conv)
  subgoal by fact
  subgoal using  $\omega_3$  unfolding  $\omega_3$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\omega_3$  unfolding  $\omega_3$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$  by auto
  subgoal using  $\omega\omega$ 
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\omega\omega$  using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
next
  case False note current = current False
  obtain w1' w2' tr1 s1' s1'' ltr1' trv1 sv1'' trv2 sv2'' statOO where
   $\chi\chi$ :  $\chi\chi$  s1 ltr1 tr1 s1' s1'' ltr1' and
   $\chi_3'$ :  $\chi_3'$   $\Delta$  w1 w2 w1' w2' s1 tr1 s1' s1'' s2 statA sv1 trv1 sv1'' sv2
trv2 sv2'' statOO
  and lltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))

  and lltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (R,w1',w2',s1'',s1'' $ ltr1',s2,ltr2,statA,sv1'',sv2'',statOO))
  using lltrv1-lltrv2-not-lnever-L[OF unw  $\Delta$  r ltr1 isi4 current]
  unfolding ltrv1 ltrv2 by blast
  define ltrv1' where ltrv1': ltrv1'  $\equiv$  lltrv1 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
  unfolding lltrv1 ltrv1' ..
  define ltrv2' where ltrv2': ltrv2'  $\equiv$  lltrv2 (R,w1',w2',s1'',s1'' $
ltr1',s2,ltr2,statA,sv1'',sv2'',statOO)
  have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
  unfolding lltrv2 ltrv2' ..

  show ?thesis apply(rule TwoFuncPred.sameFM-selectlappend)
  apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of -
w1'])

```

```

w2]) apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of -
apply(rule exI[of - ltrv1']) apply(rule exI[of - ltrv2'])
apply(intro conjI)
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrv1 .
  subgoal using ltrv2 .
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by auto
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by auto
subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by (simp add: Van.A.map-filter)

  subgoal apply(rule disjI1)
w2']) apply(rule exI[of - R]) apply(rule exI[of - w1']) apply(rule exI[of -
apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
apply(rule exI[of - s2]) apply(rule exI[of - ltr2])
apply(rule exI[of - statA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- sv2''])
apply(rule exI[of - statOO])
apply(intro conjI)
  subgoal ..subgoal ..
  subgoal unfolding ltrv1' ..
  subgoal unfolding ltrv2' ..
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def by simp
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def

  append-is-Nil-conv last-snoc not-Cons-self2 r(1))
  subgoal by fact
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
  subgoal using  $\chi^{3'}$  unfolding  $\chi^{3'}$ -def
  by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
  using llist-all-lappend-llist-of ltr1(3) by blast
  subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def using ltr2(1) by fastforce
  subgoal by fact
  subgoal by fact . .
qed
next
case R note current = current R
show ?thesis
proof(cases lnever isSecO ltr2)
  case True note current = current True
  obtain trn' w1' w2' s2' ltr2' trv1 sv1' trv2 sv2' statOO where

```

```

ωω: ltr2 = s2 $ ltr2' validTransO (s2, s2') Opt.lvalidFromS s2' ltr2'
lcompletedFromO s2' ltr2' lnever isIntO ltr2' and
ω4: ω4 Δ w1 w2 w1' w2' s1 s2 s2' statA sv1 trv1 sv1' sv2 trv2 sv2'
statOO
and trn': trn' = (if trv2 = [] then R else L)
and ltrv1: ltrv1 =
  lappend (llist-of trv1) (lltrv1 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
and ltrv2: ltrv2 =
  lappend (llist-of trv2) (lltrv2 (trn', w1', w2', s1, ltr1, s2', ltr2', statA,
sv1', sv2', statOO))
using lltrv1-lltrv2-lnever-R[OF unω Δ r ltr2(1,2) isi3 ltr2(3) current]
unfolding ltrv1 ltrv2 by blast
define ltrv1' where ltrv1': ltrv1' ≡ lltrv1 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
have lltrv1: ltrv1 = lappend (llist-of trv1) ltrv1'
unfolding ltrv1 ltrv1' ..
define ltrv2' where ltrv2': ltrv2' ≡ lltrv2 (trn', w1', w2', s1, ltr1, s2',
ltr2', statA, sv1', sv2', statOO)
have lltrv2: ltrv2 = lappend (llist-of trv2) ltrv2'
unfolding ltrv2 ltrv2' ..
show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of -
w1])
apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of -
w2])
apply(rule exI[of - ltrv1']) apply(rule exI[of - ltrv2'])
apply(intro conjI)
subgoal unfolding nL .. subgoal unfolding nR ..
subgoal using lltrv1 .
subgoal using lltrv2 .
subgoal using ω4 unfolding ω4-def by simp
subgoal using ω4 unfolding ω4-def by simp
subgoal using ω4 unfolding ω4-def by (simp add: Van.A.map-filter)
subgoal apply(rule disjI1)
apply(rule exI[of - trn']) apply(rule exI[of - w1']) apply(rule exI[of
- w2'])
apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
apply(rule exI[of - s2']) apply(rule exI[of - ltr2'])
apply(rule exI[of - statA]) apply(rule exI[of - sv1']) apply(rule exI[of
- sv2'])
apply(rule exI[of - statOO])
apply(intro conjI)
subgoal .. subgoal ..
subgoal unfolding ltrv1' ..
subgoal unfolding ltrv2' ..
subgoal using ω4 unfolding ω4-def by simp
subgoal by fact

```

```

      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Opt.reach.Step  $\omega\omega(2)$  fst-conv r(2) snd-conv)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
      subgoal using  $\omega_4$  unfolding  $\omega_4$ -def
      by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
      subgoal by fact
      subgoal by fact
      subgoal by fact
      subgoal by fact
      subgoal using  $\omega\omega$  by auto
      subgoal using  $\omega\omega$  by auto . .
next
  case False note current = current False
  obtain  $w1' w2' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statOO$  where
   $\chi\chi: \chi\chi s2 ltr2 tr2 s2' s2'' ltr2'$  and
   $\chi_4': \chi_4' \Delta w1 w2 w1' w2' s1 s2 tr2 s2' s2'' statA sv1 trv1 sv1'' sv2$ 
   $trv2 sv2'' statOO$ 
  and  $ltrv1: ltrv1 =$ 
    lappend (llist-of  $trv1$ ) (lltrv1 (L,  $w1', w2', s1, ltr1, s2'', s2'' \$ ltr2',$ 
   $statA, sv1'', sv2'', statOO$ ))
  and  $ltrv2: ltrv2 =$ 
    lappend (llist-of  $trv2$ ) (lltrv2 (L,  $w1', w2', s1, ltr1, s2'', s2'' \$ ltr2',$ 
   $statA, sv1'', sv2'', statOO$ ))
  using lltrv1-lltrv2-not-lnever-R[OF un $w \Delta r ltr2(1,2) isi3 ltr2(3)$ 
  current]
  unfolding  $ltrv1 ltrv2$  by blast
  define  $ltrv1'$  where  $ltrv1': ltrv1' \equiv lltrv1 (L, w1', w2', s1, ltr1, s2'',$ 
   $s2'' \$ ltr2', statA, sv1'', sv2'', statOO)$ 
  have lltrv1:  $ltrv1 = lappend (llist-of trv1) ltrv1'$ 
  unfolding  $ltrv1 ltrv1' ..$ 
  define  $ltrv2'$  where  $ltrv2': ltrv2' \equiv lltrv2 (L, w1', w2', s1, ltr1, s2'',$ 
   $s2'' \$ ltr2', statA, sv1'', sv2'', statOO)$ 
  have lltrv2:  $ltrv2 = lappend (llist-of trv2) ltrv2'$ 
  unfolding  $ltrv2 ltrv2' ..$ 

  show ?thesis
  apply(rule TwoFuncPred.sameFM-selectlappend)
  apply(rule exI[of -  $trv1$ ]) apply(rule exI[of -  $w1$ ]) apply(rule exI[of -
   $w1$ ])
  apply(rule exI[of -  $trv2$ ]) apply(rule exI[of -  $w2$ ]) apply(rule exI[of -
   $w2$ ])
  apply(rule exI[of -  $ltrv1$ ]) apply(rule exI[of -  $ltrv2$ ])
  apply(intro conjI)
  subgoal unfolding  $nL ..$  subgoal unfolding  $nR ..$ 
  subgoal using  $lltrv1 .$ 
  subgoal using  $lltrv2 .$ 

```

```

      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
      subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
    subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by (simp add: Van.A.map-filter)
      subgoal apply(rule disjI1)
        apply(rule exI[of - L]) apply(rule exI[of - w1']) apply(rule exI[of -
w2'])
          apply(rule exI[of - s1]) apply(rule exI[of - ltr1])
          apply(rule exI[of - s2']) apply(rule exI[of - s2'' $ ltr2'])
          apply(rule exI[of - statA])
        apply(rule exI[of - sv1']) apply(rule exI[of - sv2']) apply(rule exI[of
- statOO])
          apply(intro conjI)
            subgoal .. subgoal ..
              subgoal unfolding ltrv1' ..
              subgoal unfolding ltrv2' ..
              subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def by simp
              subgoal by fact
              subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
            by (metis Simple-Transition-System.reach-validFromS-reach  $\chi\chi$   $\chi\chi$ -def
              append-is-Nil-conv last-snoc not-Cons-self2 r(2))
          subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
            by (metis Simple-Transition-System.reach-validFromS-reach r(3)
snoc-eq-iff-butlast)
          subgoal using  $\chi_4'$  unfolding  $\chi_4'$ -def
            by (metis Simple-Transition-System.reach-validFromS-reach r(4)
snoc-eq-iff-butlast)
          subgoal by fact
          subgoal by fact
          subgoal by fact
          subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
          subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def by auto
          subgoal using  $\chi\chi$  unfolding  $\chi\chi$ -def
          using llist-all-lappend-llist-of ltr2(3) by blast . .
        qed
      qed
    qed
  qed
}
}
thus ?thesis unfolding Van.lA[OF cltrv1] Van.lA[OF cltrv2]
  apply- apply(rule TwoFuncPred.sameFM-lmap-lfilter)
  using assms by blast
qed

```



```

fun isN :: ('stateO,'stateV) tuple34 ⇒ bool
where
isN (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ⇔ ltr1 = [] ∨ ltr2 = []

fun H-T ::
('stateO,'stateV)tuple34 ⇒
('stateO,'stateV)tuple12 ×
('stateO,'stateV)tuple34
where
H-T (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) =
  (let (tr1,s1',s1'',ltr1',tr2,s2',s2'',ltr2') =
      (SOME k. case k of (tr1,s1',s1'',ltr1',tr2,s2',s2'',ltr2') ⇒
        φφ s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2')
    in let (w1',w2',trv1,sv1'',trv2,sv2'',statAA,statOO) =
        (SOME k'. case k' of (w1',w2',trv1,sv1'',trv2,sv2'',statAA,statOO) ⇒
          φ' Δ w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA statO
        sv1 trv1 sv1'' sv2 trv2 sv2'' statOO)
      in ((trv1,sv1'',trv2,sv2'',statAA,statOO),
        (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))
    )

declare H-T.simps[simp del]

definition H ≡ fst o H-T
definition T ≡ snd o H-T

fun Econd where Econd (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) = (lnever
isIntO ltr1)

fun E where E (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) = f (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)

definition F :: ('stateO,'stateV)tuple34 ⇒ ('stateO,'stateV)tuple12 llist
where F ≡ ccorec-llist isN H Econd E T

```

```

lemma F-LNil:
ltr1 = [] ∨ ltr2 = [] ⇒ F (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) = []
unfolding F-def apply(subst llist-ccorec(1)) by auto

```

```

lemma F-lnever:
assumes ltr1 ≠ [] ltr2 ≠ [] lnever isIntO ltr1
shows F (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) = f (L, w1, w2, s1, ltr1, s2,
ltr2, statA, sv1, sv2, statO)
using assms unfolding F-def apply(subst llist-ccorec(2))
  subgoal unfolding E.simps lnull-def apply(rule f-not-LNil) by auto
  subgoal using assms by auto
  subgoal unfolding Econd.simps by auto .

```

lemma *F-not-lnever*:

assumes $ltr1 \neq []$ $ltr2 \neq []$ $\neg lnever$ *isIntO* $ltr1$

shows $F (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$

$LCons (H (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)) (F (T (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)))$

using *assms* **unfolding** *F-def* **apply**(*subst llist-ccorec(2)*)

subgoal unfolding *E.simps lnull-def* **apply**(*rule f-not-LNil*) **by auto**

subgoal using *assms* **by auto**

subgoal unfolding *Econd.simps* **by auto** .

definition *ltrv1* :: $('stateO, 'stateV)tuple34 \Rightarrow 'stateV$ **l***list* **where**

$ltrv1\ tp = lconcat (lmap (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). llist-of\ trv1) (F\ tp))$

definition *firstHolds1* :: $('stateO, 'stateV)tuple34 \Rightarrow nat$ **where**

$firstHolds1\ tp = firstNC (lmap (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). trv1) (F\ tp))$

definition *ltrv2* :: $('stateO, 'stateV)tuple34 \Rightarrow 'stateV$ **l***list* **where**

$ltrv2\ tp = lconcat (lmap (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). llist-of\ trv2) (F\ tp))$

definition *firstHolds2* :: $('stateO, 'stateV)tuple34 \Rightarrow nat$ **where**

$firstHolds2\ tp = firstNC (lmap (\lambda(trv1, sv1'', trv2, sv2'', statAA, statOO). trv2) (F\ tp))$

lemma *ltrv1-ne-imp*:

assumes $ltrv1\ tp \neq []$

shows $\exists trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO. (trv1, sv1'', trv2, sv2'', statAA, statOO) \in lset (F\ tp) \wedge$

$trv1 \neq []$

using *assms* **unfolding** *ltrv1-def* **unfolding** *lconcat-eq-LNil-iff* **by force**

lemma *ltrv2-ne-imp*:

assumes $ltrv2\ tp \neq []$

shows $\exists trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO. (trv1, sv1'', trv2, sv2'', statAA, statOO) \in lset (F\ tp) \wedge$

$trv2 \neq []$

using *assms* **unfolding** *ltrv2-def* **unfolding** *lconcat-eq-LNil-iff* **by force**

lemma *ltrv1-LNil[simp]*:

$ltr1 = [] \vee ltr2 = [] \implies ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = []$
unfolding $ltrv1\text{-def } F\text{-LNil}$ **by** $simp$
lemma $ltrv2\text{-LNil}[simp]$:
 $ltr1 = [] \vee ltr2 = [] \implies ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = []$
unfolding $ltrv2\text{-def } F\text{-LNil}$ **by** $simp$

lemma $ltrv1\text{-lnever}$:
assumes $ltr1 \neq []$ $ltr2 \neq []$ $lnever\ isIntO\ ltr1$
shows $ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$
unfolding $ltrv1\text{-def } F\text{-lnever}[OF\ assms]$ $lltrv1\text{-def}$ **..**

lemma $ltrv2\text{-lnever}$:
assumes $ltr1 \neq []$ $ltr2 \neq []$ $lnever\ isIntO\ ltr1$
shows $ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)$
unfolding $ltrv2\text{-def } F\text{-lnever}[OF\ assms]$ $lltrv2\text{-def}$ **..**

lemma $H\text{-T-not-lnever}$:
assumes $unw: unwindCond\ \Delta$
and $\Delta: \Delta\ \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and $r: reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$
and $stat: statA = Diff \longrightarrow statO = Diff$
and $ltr1: Opt.lvalidFromS\ s1\ ltr1\ Opt.lcompletedFrom\ s1\ ltr1$
and $ltr2: Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$
and $l: \neg\ lnever\ isIntO\ ltr1\ Opt.lA\ ltr1 = Opt.lA\ ltr2$
shows $\exists\ w1'\ w2'\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'\ trv1\ sv1''\ trv2\ sv2''\ statAA\ statOO.$

$\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2' \wedge$
 $\varphi'\ \Delta\ w1\ w2\ w1'\ w2'\ statA\ s1\ tr1\ s1'\ s1''\ s2\ tr2\ s2'\ s2''\ statAA\ statO\ sv1\ trv1$
 $sv1''\ sv2\ trv2\ sv2''\ statOO \wedge$
 $H\text{-T}\ (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $((trv1, sv1'', trv2, sv2'', statAA, statOO),$
 $(w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$

proof –
obtain $tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$
where $\varphi\varphi: \varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2'$
using $isIntO\text{-}\varphi\varphi[OF\ ltr1\ ltr2\ l(2,1)]$
by $auto$

define tp **where**
 $tp = (SOME\ k.\ case\ k\ of\ (tr1, s1', s1'', ltr1', tr2, s2', s2'', ltr2') \Rightarrow$
 $\varphi\varphi\ s1\ ltr1\ s2\ ltr2\ tr1\ s1'\ s1''\ ltr1'\ tr2\ s2'\ s2''\ ltr2')$

have $0: case\ tp\ of\ (tr1, s1', s1'', ltr1', tr2, s2', s2'', ltr2') \Rightarrow$

```

       $\varphi\varphi$   $s1$   $ltr1$   $s2$   $ltr2$   $tr1$   $s1'$   $s1''$   $ltr1'$   $tr2$   $s2'$   $s2''$   $ltr2'$ 
using  $\varphi\varphi$  unfolding  $tp$ -def apply- apply(rule someI-ex)
apply(rule exI[of - (tr1,s1',s1'',ltr1',tr2,s2',s2'',ltr2')]) by auto

obtain  $tr1$   $s1'$   $s1''$   $ltr1'$   $tr2$   $s2'$   $s2''$   $ltr2'$  where
 $tp$ :  $tp = (tr1,s1',s1'',ltr1',tr2,s2',s2'',ltr2')$  by(cases  $tp$ , auto)

have  $\varphi\varphi$ :  $\varphi\varphi$   $s1$   $ltr1$   $s2$   $ltr2$   $tr1$   $s1'$   $s1''$   $ltr1'$   $tr2$   $s2'$   $s2''$   $ltr2'$ 
using 0 unfolding  $tp$  by auto

obtain  $w1'$   $w2'$   $trv1$   $sv1''$   $trv2$   $sv2''$   $statAA$   $statOO$ 
where  $\varphi'$ :  $\varphi' \Delta w1$   $w2$   $w1'$   $w2'$   $statA$   $s1$   $tr1$   $s1'$   $s1''$   $s2$   $tr2$   $s2'$   $s2''$   $statAA$   $statO$ 
 $sv1$   $trv1$   $sv1''$   $sv2$   $trv2$   $sv2''$   $statOO$ 
using unwindCond-ex- $\varphi'$ [OF un $w$   $\Delta$   $r$   $stat$ , of  $tr1$   $s1'$   $s1''$   $tr2$   $s2'$   $s2''$ ]
using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by auto

define  $tp'$  where
 $tp' = (SOME$   $k'$ . case  $k'$  of ( $w1'$ , $w2'$ , $trv1$ , $sv1''$ , $trv2$ , $sv2''$ , $statAA$ , $statOO$ )  $\Rightarrow$ 
 $\varphi' \Delta w1$   $w2$   $w1'$   $w2'$   $statA$   $s1$   $tr1$   $s1'$   $s1''$   $s2$   $tr2$   $s2'$   $s2''$   $statAA$   $statO$   $sv1$ 
 $trv1$   $sv1''$   $sv2$   $trv2$   $sv2''$   $statOO$ )

have 1: case  $tp'$  of ( $w1'$ , $w2'$ , $trv1$ , $sv1''$ , $trv2$ , $sv2''$ , $statAA$ , $statOO$ )  $\Rightarrow$ 
 $\varphi' \Delta w1$   $w2$   $w1'$   $w2'$   $statA$   $s1$   $tr1$   $s1'$   $s1''$   $s2$   $tr2$   $s2'$   $s2''$   $statAA$   $statO$   $sv1$ 
 $trv1$   $sv1''$   $sv2$   $trv2$   $sv2''$   $statOO$ 
using  $\varphi'$  unfolding  $tp'$ -def apply- apply(rule someI-ex)
apply(rule exI[of - ( $w1'$ , $w2'$ , $trv1$ , $sv1''$ , $trv2$ , $sv2''$ , $statAA$ , $statOO$ )]) by auto

obtain  $w1'$   $w2'$   $trv1$   $sv1''$   $trv2$   $sv2''$   $statAA$   $statOO$  where
 $tp'$ :  $tp' = (w1',w2',trv1,sv1'',trv2,sv2'',statAA,statOO)$  by(cases  $tp'$ , auto)

have  $\varphi'$ :  $\varphi' \Delta w1$   $w2$   $w1'$   $w2'$   $statA$   $s1$   $tr1$   $s1'$   $s1''$   $s2$   $tr2$   $s2'$   $s2''$   $statAA$   $statO$ 
 $sv1$   $trv1$   $sv1''$   $sv2$   $trv2$   $sv2''$   $statOO$ 
using 1 unfolding  $tp'$  by auto

show ?thesis
apply(rule exI[of -  $w1'$ ]) apply(rule exI[of -  $w2'$ ])
apply(rule exI[of -  $tr1$ ]) apply(rule exI[of -  $s1'$ ]) apply(rule exI[of -  $s1''$ ])
apply(rule exI[of -  $ltr1'$ ])
apply(rule exI[of -  $tr2$ ]) apply(rule exI[of -  $s2'$ ]) apply(rule exI[of -  $s2''$ ])
apply(rule exI[of -  $ltr2'$ ])
apply(rule exI[of -  $trv1$ ]) apply(rule exI[of -  $sv1''$ ]) apply(rule exI[of -  $trv2$ ])
apply(rule exI[of -  $sv2''$ ])
apply(rule exI[of -  $statAA$ ]) apply(rule exI[of -  $statOO$ ])
apply(intro conjI)
subgoal using  $\varphi\varphi$  .
subgoal using  $\varphi'$  .
subgoal unfolding H-T.simps

```

unfolding $tp\text{-def}[symmetric]$ tp **apply** $simp$
unfolding $tp'\text{-def}[symmetric]$ tp' **by** $simp$.
qed

lemma $ltrv1\text{-}ltrv2\text{-not}\text{-lnever}$:

assumes unw : $unwindCond$ Δ

and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$

and r : $reachO s1 reachO s2 reachV sv1 reachV sv2$

and $stat$: $statA = Diff \longrightarrow statO = Diff$

and $ltr1$: $Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1$

and $ltr2$: $Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2$

and l : $\neg lnever isIntO ltr1 Opt.lA ltr1 = Opt.lA ltr2$

shows $\exists w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA statOO$.

$\varphi\varphi s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' \wedge$
 $\varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA statO sv1 trv1$
 $sv1'' sv2 trv2 sv2'' statOO \wedge$
 $ltrv1 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend (l\text{-list-of } trv1) (ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$
 \wedge
 $ltrv2 (w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO) =$
 $lappend (l\text{-list-of } trv2) (ltrv2 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$

proof–

have $ltr1NE$: $ltr1 \neq []$ **using** $l(1)$ **by** $auto$

hence $ltr2NE$: $ltr2 \neq []$ **using** $l(2)$

using $Opt.lcompletedFrom\text{-}def$ $ltr2(2)$ **by** $blast$

show $?thesis$

using $H\text{-}T\text{-not}\text{-lnever}[OF assms]$ **apply** $(elim exE)$

subgoal for $w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA statOO$

apply $(rule exI[of - w1'])$ **apply** $(rule exI[of - w2'])$

apply $(rule exI[of - tr1])$ **apply** $(rule exI[of - s1'])$ **apply** $(rule exI[of - s1''])$

apply $(rule exI[of - ltr1'])$

apply $(rule exI[of - tr2])$ **apply** $(rule exI[of - s2'])$ **apply** $(rule exI[of - s2''])$

apply $(rule exI[of - ltr2'])$

apply $(rule exI[of - trv1])$ **apply** $(rule exI[of - sv1''])$ **apply** $(rule exI[of - trv2])$

apply $(rule exI[of - sv2''])$

apply $(rule exI[of - statAA])$ **apply** $(rule exI[of - statOO])$

apply $(intro conjI)$

subgoal by $simp$

subgoal by $simp$

subgoal unfolding $ltrv1\text{-}def$ **apply** $(subst F\text{-not}\text{-lnever}[OF ltr1NE ltr2NE l(1)])$

unfolding $H\text{-}def$ $T\text{-}def$ **by** $simp$

subgoal unfolding $ltrv2\text{-}def$ **apply** $(subst F\text{-not}\text{-lnever}[OF ltr1NE ltr2NE l(1)])$

unfolding $H\text{-}def$ $T\text{-}def$ **by** $simp$. .

qed

lemma *lvalidFromS-ltrv1*:
assumes *unw*: *unwindCond* Δ
and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *stat*: *statA* = *Diff* \longrightarrow *statO* = *Diff*
and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1*
and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2*
and *ltr14*: *Opt.lA* *ltr1* = *Opt.lA* *ltr2*
shows *Van.lvalidFromS* *sv1* (*ltrv1* (*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*))
proof –
 {**fix** *n1* *sv1* *ltrr1*
 assume $\exists w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv2\ statO.$
 ltrr1 = *ltrv1* (*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*) \wedge
 n1 = *w1* \wedge
 $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 reachO *s1* \wedge *reachO* *s2* \wedge *reachV* *sv1* \wedge *reachV* *sv2* \wedge
 (*statA* = *Diff* \longrightarrow *statO* = *Diff*) \wedge
 Opt.lvalidFromS *s1* *ltr1* \wedge *Opt.lcompletedFrom* *s1* *ltr1* \wedge
 Opt.lvalidFromS *s2* *ltr2* \wedge *Opt.lcompletedFrom* *s2* *ltr2* \wedge
 Opt.lA *ltr1* = *Opt.lA* *ltr2*
 hence *Van.llvalidFromS* *n1* *sv1* *ltrr1*
 proof(*coinduct* rule: *Van.llvalidFromS.coinduct*[of $\lambda n1\ sv1\ ltrr1.$
 $\exists w1\ w2\ s1\ ltr1\ s2\ ltr2\ statA\ sv2\ statO.$
 ltrr1 = *ltrv1* (*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*) \wedge
 n1 = *w1* \wedge
 $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 reachO *s1* \wedge *reachO* *s2* \wedge *reachV* *sv1* \wedge *reachV* *sv2* \wedge
 (*statA* = *Diff* \longrightarrow *statO* = *Diff*) \wedge
 Opt.lvalidFromS *s1* *ltr1* \wedge *Opt.lcompletedFrom* *s1* *ltr1* \wedge
 Opt.lvalidFromS *s2* *ltr2* \wedge *Opt.lcompletedFrom* *s2* *ltr2* \wedge
 Opt.lA *ltr1* = *Opt.lA* *ltr2*])
 case (*llvalidFromS* *n1* *sv1* *ltrr1*)
 then obtain *w1* *w2* *s1* *ltr1* *s2* *ltr2* *statA* *sv2* *statO*
 where *ltrr1*: *ltrr1* = *ltrv1* (*w1*,*w2*,*s1*,*ltr1*,*s2*,*ltr2*,*statA*,*sv1*,*sv2*,*statO*)
 and *n1*: *n1* = *w1*
 and Δ : $\Delta \infty w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
 and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
 and *stat*: *statA* = *Diff* \longrightarrow *statO* = *Diff*
 and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1*
 and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2*
 and *A34*: *Opt.lA* *ltr1* = *Opt.lA* *ltr2*
 by *auto*

 have *current*: *ltr1* \neq $\llbracket \ \ \rrbracket$ *ltr2* \neq $\llbracket \ \ \rrbracket$
 using *ltr1*(2) *ltr2*(2) **unfolding** *Opt.lcompletedFrom-def* **by** *auto*

 show ?*case proof*(*cases* *lnever isIntO* *ltr1*)
 case *True* **note** *current* = *current* *True*
 hence *lnever isIntO* *ltr2*

```

    by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
    note ln34 = True this
    have ltrr1: ltrr1 = ltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
    unfolding ltrr1 ltrv1-lnever[OF current] by simp
    show ?thesis apply(rule Van.lvalidFromS-selectlvalidFromS)
    unfolding ltrr1 apply simp
    apply(rule lvalidFromS-ltrv1)
    using ln34  $\Delta$  lvalidFromS ln34(2) ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2)
r(4) unw by auto
  next
  case False note ln3 = False
  hence ln4:  $\neg$  lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr1  $\neq$  [[s1]] using ln3 ltr1
  using Opt.final-not-isInt by auto
  hence llength ltr1 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
  hence  $\neg$  finalO s1
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
  hence nf12:  $\neg$  finalV sv1  $\wedge$   $\neg$  finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unw unwindCond-def by force

  obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and  $\varphi'$ :  $\varphi'$   $\Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and ltrr1: ltrr1 =
lappend (llist-of trv1) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

  using ltrv1-ltrv2-not-lnever[OF unw  $\Delta$  r stat ltr1 ltr2 ln3 A34]
  unfolding ltrr1 by blast
  define ltrr1' where ltrr1': ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
  unfolding ltrr1 ltrr1' ..
  have ne: trv1  $\neq$  []  $\vee$  (trv1 = []  $\wedge$  w1' < w1)
  using  $\varphi'$  unfolding  $\varphi'$ -def ltrr1 by simp

  show ?thesis using ne proof(elim disjE conjE)
    assume trv1: trv1  $\neq$  []
    show ?thesis

```

```

apply(rule Van.llvalidFromS-selectlappend)
apply(rule exI[of - sv1]) apply(rule exI[of - trv1])
apply(rule exI[of - sv1'']) apply(rule exI[of - w1'])
apply(rule exI[of - lrr1']) apply(rule exI[of - w1])
apply(intro conjI)
  subgoal unfolding n1 .. subgoal ..
  subgoal unfolding lrr1 ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.validS-append1 Van.validFromS-def append-is-Nil-conv
hd-append2)
  subgoal by fact
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def
  by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
append-is-Nil-conv list.sel(1) not-Cons-self2 trv1)
  subgoal apply(rule disjI1)
  apply(rule exI[of - w1']) apply(rule exI[of - w2'])
  apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ lrr1'])
  apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ lrr2'])
  apply(rule exI[of - statAA]) apply(rule exI[of - sv2'']) apply(rule exI[of
- statOO])
  apply(intro conjI)
  subgoal unfolding lrr1' ..
  subgoal ..
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach append-is-Nil-conv last-snoc
trv1)
  subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
next
assume trv1[simp]: trv1 = [] and MM': w1' < w1
hence sv1''[simp]: sv1'' = sv1 using  $\varphi'$  unfolding  $\varphi'$ -def by simp
show ?thesis
apply(rule Van.llvalidFromS-selectDelay)
apply(rule exI[of - w1']) apply(rule exI[of - w1])
apply(rule exI[of - sv1'']) apply(rule exI[of - lrr1'])
apply(intro conjI)

```



```

subgoal unfolding n1 .. subgoal by simp
subgoal unfolding ltrr1 by simp subgoal by fact
subgoal apply(rule disjI1)
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
apply(rule exI[of - statAA]) apply(rule exI[of - sv2'']) apply(rule exI[of
- statOO])
apply(intro conjI)
subgoal unfolding ltrr1' ..
subgoal ..
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
subgoal unfolding sv1'' by fact
subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
qed
qed
qed
}
thus ?thesis apply-apply(rule Van.lvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

```

lemma lvalidFromS-ltrv2:
assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and ltr14: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lvalidFromS sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof -
{fix n2 sv2 ltrr2
assume  $\exists w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO.$ 
ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
n2 = w2  $\wedge$ 

```

```

    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    (statA = Diff → statO = Diff) ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
    Opt.lA ltr1 = Opt.lA ltr2
  hence Van.llvalidFromS n2 sv2 ltr2
  proof (coinduct rule: Van.llvalidFromS.coinduct[of λn2 sv2 ltr2.
    ∃ w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO.
      ltr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
      n2 = w2 ∧
      Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
      reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
      (statA = Diff → statO = Diff) ∧
      Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
      Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
      Opt.lA ltr1 = Opt.lA ltr2])
  case (llvalidFromS n2 sv2 ltr2)
  then obtain w1 w2 s1 ltr1 s2 ltr2 statA sv1 statO
  where ltr2: ltr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and n2: n2 = w2
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff → statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  by auto

  have current: ltr1 ≠ [] ltr2 ≠ []
  using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

  show ?case proof (cases lnever isIntO ltr1)
  case True note current = current True
  hence lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
    ltr2(2))
  note ln34 = True this
  have ltr2: ltr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
    statO)
  unfolding ltr2 ltrv2-lnever[OF current] by simp
  show ?thesis apply (rule Van.llvalidFromS-selectlvalidFromS)
  unfolding ltr2 apply simp
  apply (rule lvalidFromS-lltrv2)
  using ln34 Δ llvalidFromS ln34(2) ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2)
  r(3) unω by auto
  next
  case False note ln3 = False
  hence ln4: ¬ lnever isIntO ltr2

```

by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

have ltr2 ≠ [[s2]] using ln4 ltr2
 using Opt.final-not-isInt by auto
 hence llength ltr2 > Suc 0
 by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(2) enat-0-iff(2)
 linorder-not-less llength-LNil llist-eq-cong ltr2(1) ltr2(2) nle-le not-iless0)
 hence ¬ finalO s2
 by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(2) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr2(1))
 hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
 using Δ r(1) r(2) r(3) r(4) unw unwindCond-def by force

obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
 statOO

where φφ: φφ s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
 and φ': φ' Δ w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
 statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
 and ltrr2: ltrr2 =
 lappend (llist-of trv2) (ltrv2 (w1',w2',s1'',s1'' \$ ltr1',s2'',s2'' \$ ltr2',statAA,sv1'',sv2'',statOO))

using ltrv1-ltrv2-not-lnever[OF unw Δ r stat ltr1 ltr2 ln3 A34]

unfolding ltrr2 by blast

define ltrr2' where ltrr2': ltrr2' = ltrv2 (w1',w2',s1'',s1'' \$ ltr1',s2'',s2''
 \$ ltr2',statAA,sv1'',sv2'',statOO)

have ltrr2: ltrr2 = lappend (llist-of trv2) ltrr2'

unfolding ltrr2 ltrr2' ..

have ne: trv2 ≠ [] ∨ (trv2 = [] ∧ w2' < w2)

using φ' unfolding φ'-def ltrr2 by simp

show ?thesis using ne proof(elim disjE conjE)

assume trv2: trv2 ≠ []

show ?thesis

apply(rule Van.lvalidFromS-selectlappend)

apply(rule exI[of - sv2]) apply(rule exI[of - trv2])

apply(rule exI[of - sv2'']) apply(rule exI[of - w2'])

apply(rule exI[of - ltrr2']) apply(rule exI[of - w2])

apply(intro conjI)

subgoal unfolding n2 .. subgoal ..

subgoal unfolding ltrr2 ..

subgoal using φ' unfolding φ'-def

by (metis Van.validS-append1 Van.validFromS-def append-is-Nil-conv
 hd-append2)

subgoal by fact

subgoal using φ' unfolding φ'-def

by (metis Simple-Transition-System.validFromS-def Van.validS-validTrans
 append-is-Nil-conv list.distinct(1) list.sel(1) trv2)

```

subgoal apply(rule disjI1)
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
apply(rule exI[of - statAA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- statOO])
apply(intro conjI)
subgoal unfolding ltrr2' ..
subgoal ..
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
next
assume trv2[simp]: trv2 = [] and MM': w2' < w2
hence sv2''[simp]: sv2'' = sv2 using  $\varphi'$  unfolding  $\varphi'$ -def by simp
show ?thesis
apply(rule Van.lvalidFromS-selectDelay)
apply(rule exI[of - w2']) apply(rule exI[of - w2])
apply(rule exI[of - sv2'']) apply(rule exI[of - ltrr2'])
apply(intro conjI)
subgoal unfolding n2 .. subgoal by simp
subgoal unfolding ltrr2 by simp subgoal by fact
subgoal apply(rule disjI1)
apply(rule exI[of - w1']) apply(rule exI[of - w2'])
apply(rule exI[of - s1'']) apply(rule exI[of - s1'' $ ltr1'])
apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
apply(rule exI[of - statAA]) apply(rule exI[of - sv1'']) apply(rule exI[of
- statOO])
apply(intro conjI)
subgoal unfolding ltrr2' ..
subgoal ..
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)

```

```

      subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
        by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
      subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
        by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
      subgoal unfolding sv2'' by fact
      subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
      subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
    qed
  qed
  qed
}
thus ?thesis apply-apply(rule Van.lvalidFromS-imp-lvalidFromS)
using assms by blast
qed

```

```

lemma lcompletedFrom-ltrv1:
  assumes unw: unwindCond  $\Delta$ 
  and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff  $\longrightarrow$  statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  shows Van.lcompletedFrom sv1 (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  proof-
    {fix ltrr1 assume ltrr1: ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
      and lfin: lfinite ltrr1
      hence list-of ltrr1  $\neq [] \wedge$  finalV (last (list-of ltrr1))
      using assms(2-) proof(induct length (list-of ltrr1) w1
        arbitrary: w2 ltrr1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
        rule: less2-induct')
        case (less w1 ltrr1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)
          hence ltrr1: ltrr1 = ltrv1 (w1,w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)
          and lfin: lfinite ltrr1
          and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
          and r: reachO s1 reachO s2 reachV sv1 reachV sv2
          and stat: statA = Diff  $\longrightarrow$  statO = Diff
          and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1
          and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2
          and A34: Opt.lA ltr1 = Opt.lA ltr2
          by auto
        }
  }

```

```

have current: ltr1 ≠ [[]] ltr2 ≠ [[]]
using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

show ?case proof(cases lnever isIntO ltr1)
  case True note ln3 = True note current = current True
  hence ln4: lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
  note ln34 = True this
  have ltrr1: ltrr1 = ltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr1 ltrv1-lnever[OF current] by simp
  show ?thesis
  using lcompletedFrom-ltrv1[OF unω Δ r ltr1 ln3 ltr2 ln4, of L]
  using lfin[unfolded ltrr1]
  unfolding Van.lcompletedFrom-def ltrr1[symmetric]
  using llist-of-list-of by fastforce
next
  case False note ln3 = False
  hence ln4: ¬ lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr1 ≠ [[s1]] using ln3 ltr1
  using Opt.final-not-isInt by auto
  hence llength ltr1 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
  hence ¬ finalO s1
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
  hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unω unwindCond-def by force

  obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and  $\varphi'$ :  $\varphi'$   $\Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and ltrr1: ltrr1 =
lappend (llist-of trv1) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

  using ltrv1-ltrv2-not-lnever[OF unω Δ r stat ltr1 ltr2 ln3 A34]
  unfolding ltrr1 by blast
  define ltrr1' where ltrr1': ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
 $\$$  ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
  unfolding ltrr1 ltrr1' ..

```

```

have ne: trv1 ≠ [] ∨ (trv1 = [] ∧ w1' < w1)
using  $\varphi'$  unfolding  $\varphi'$ -def lrr1 by simp

have lfin': lfinite lrr1'
using lfin ne unfolding lrr1 by simp
have len: length (list-of lrr1') < length (list-of lrr1) ∨
      length (list-of lrr1') = length (list-of lrr1) ∧ w1' < w1
using ne lfin lfin' by (simp add: list-of-lappend lrr1)

have 0: list-of lrr1' ≠ [] ∧ finalV (last (list-of lrr1'))
using len proof(elim disjE conjE)
  assume len: length (list-of lrr1') < length (list-of lrr1)
  show ?thesis
  apply(rule less(1)[OF - lrr1'])
    subgoal by fact subgoal by fact
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp .
next
  assume len: length (list-of lrr1') = length (list-of lrr1) w1' < w1
  show ?thesis
  apply(rule less(2)[OF - - lrr1'])
    subgoal by fact subgoal unfolding len ..
    subgoal by fact
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)

```

```

    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by auto
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp .
  qed
  show ?thesis unfolding ltrr1 using 0
  by (simp add: lfin' list-of-lappend)
  qed
  qed
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

```

```

lemma lcompletedFrom-ltrv2:
  assumes unvw: unwindCond  $\Delta$ 
  and  $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff  $\longrightarrow$  statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  shows Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  proof -
    {fix ltrr2 assume ltrr2: ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
      and lfin: lfinite ltrr2
      hence list-of ltrr2  $\neq$  []  $\wedge$  finalV (last (list-of ltrr2))
      using assms(2-) proof(induct length (list-of ltrr2) w2
        arbitrary: w1 ltrr2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
        rule: less2-induct')
        case (less w2 ltrr2 w1 s1 ltr1 s2 ltr2 statA sv1 sv2 statO)
          hence ltrr2: ltrr2 = ltrv2 (w1,w2, s1, ltr1, s2, ltr2, statA, sv1, sv2, statO)
          and lfin: lfinite ltrr2
          and  $\Delta: \Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
          and r: reachO s1 reachO s2 reachV sv1 reachV sv2
          and stat: statA = Diff  $\longrightarrow$  statO = Diff
          and ltr1: Opt.lvalidFromS s1 ltr1 lcompletedFromO s1 ltr1
          and ltr2: Opt.lvalidFromS s2 ltr2 lcompletedFromO s2 ltr2
          and A34: Opt.lA ltr1 = Opt.lA ltr2
          by auto
        }
    have current: ltr1  $\neq$  [] ltr2  $\neq$  []
    using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

  show ?case proof(cases lnever isIntO ltr1)
    case True note ln3 = True note current = current True
    hence ln4: lnever isIntO ltr2
    by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34)
  
```



```

ltr2(2))
  note ln34 = True this
  have ltrr2: ltrr2 = ltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr2 ltrv2-lnever[OF current] by simp
  show ?thesis
  using lcompletedFrom-ltrv2[OF unW Δ r ltr1 ln3 ltr2 ln4, of L]
  using lfin[unfolded ltrr2]
  unfolding Van.lcompletedFrom-def ltrr2[symmetric]
  using llist-of-list-of by fastforce
next
case False note ln3 = False
hence ln4: ¬ lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr2 ≠ [[s2]] using ln4 ltr2
  using Opt.final-not-isInt by auto
  hence llength ltr2 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(2) enat-0-iff(2)
linorder-not-less llength-LNil llist-eq-cong ltr2(1) ltr2(2) nle-le not-less-zero)
  hence ¬ finalO s2
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(2) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr2(1))
  hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
  using Δ r(1) r(2) r(3) r(4) unW unwindCond-def by force

  obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where φφ: φφ s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and φ': φ' Δ w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
  and ltrr2: ltrr2 =
lappend (llist-of trv2) (ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

  using ltrv1-ltrv2-not-lnever[OF unW Δ r stat ltr1 ltr2 ln3 A34]
  unfolding ltrr2 by blast
  define ltrr2' where ltrr2': ltrr2' = ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
  have ltrr2: ltrr2 = lappend (llist-of trv2) ltrr2'
  unfolding ltrr2 ltrr2' ..
  have ne: trv2 ≠ [] ∨ (trv2 = [] ∧ w2' < w2)
  using φ' unfolding φ'-def ltrr2 by simp

  have lfin': lfinite ltrr2'
  using lfin ne unfolding ltrr2 by simp
  have len: length (list-of ltrr2') < length (list-of ltrr2) ∨
length (list-of ltrr2') = length (list-of ltrr2) ∧ w2' < w2

```

```

using ne lfin lfin' by (simp add: list-of-lappend ltrr2)

have 0: list-of ltrr2' ≠ [] ∧ finalV (last (list-of ltrr2'))
using len proof (elim disjE conjE)
  assume len: length (list-of ltrr2') < length (list-of ltrr2)
  show ?thesis
  apply (rule less(1)[OF - ltrr2'])
    subgoal by fact subgoal by fact
    subgoal using φ' unfolding φ'-def by simp
    subgoal using r(1) φφ unfolding φφ-def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(2) φφ unfolding φφ-def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3) φ' unfolding φ'-def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using r(4) φ' unfolding φ'-def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using φ' unfolding φ'-def by auto
    subgoal using φφ unfolding φφ-def by simp
    subgoal using φφ unfolding φφ-def by simp
    subgoal using φφ unfolding φφ-def by simp
    subgoal using φφ unfolding φφ-def by simp
    subgoal using φφ unfolding φφ-def by simp .
next
assume len: length (list-of ltrr2') = length (list-of ltrr2) w2' < w2
show ?thesis
apply (rule less(2)[OF - - ltrr2'])
  subgoal by fact subgoal unfolding len ..
  subgoal by fact
  subgoal using φ' unfolding φ'-def by simp
  subgoal using r(1) φφ unfolding φφ-def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(2) φφ unfolding φφ-def
    by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
  subgoal using r(3) φ' unfolding φ'-def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using r(4) φ' unfolding φ'-def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
  subgoal using φ' unfolding φ'-def by auto
  subgoal using φφ unfolding φφ-def by simp
  subgoal using φφ unfolding φφ-def by simp
  subgoal using φφ unfolding φφ-def by simp
  subgoal using φφ unfolding φφ-def by simp
  subgoal using φφ unfolding φφ-def by simp .
qed

```

```

    show ?thesis unfolding ltrr2 using 0
    by (simp add: lfin' list-of-lappend)
  qed
qed
}
thus ?thesis unfolding Van.lcompletedFrom-def by auto
qed

```

lemma *lS-ltrv1-ltr1*:

```

assumes unw: unwindCond Δ
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff ⟶ statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lS (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) = Opt.lS ltr1
proof-
  have cltrv1: Van.lcompletedFrom sv1 (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-ltrv1[OF assms] .
  {fix nL nR ltrr1 ltr1
  assume ∃ w1 w2 s1 s2 ltr2 statA sv1 sv2 statO.
    nL = w1 ∧ nR = w1 ∧
    ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    (statA = Diff ⟶ statO = Diff) ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
    Opt.lA ltr1 = Opt.lA ltr2
  hence TwoFuncPred.sameFM isSecV isSecO getSecV getSecO nL nR ltrr1 ltr1
  proof(coinduct rule: TwoFuncPred.sameFM.coinduct[of λnL nR ltrr1 ltr1.
    ∃ w1 w2 s1 s2 ltr2 statA sv1 sv2 statO.
      nL = w1 ∧ nR = w1 ∧
      ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
      Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
      reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
      (statA = Diff ⟶ statO = Diff) ∧
      Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
      Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
      Opt.lA ltr1 = Opt.lA ltr2])
    case (2 nL nR ltrr1 ltr1)
    then obtain w1 w2 s1 s2 ltr2 statA sv1 sv2 statO
    where nL: nL = w1 and nR: nR = w1
    and ltrr1: ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
    and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
    and r: reachO s1 reachO s2 reachV sv1 reachV sv2

```

```

and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
by auto

have current: ltr1  $\neq$   $[\ ]$  ltr2  $\neq$   $[\ ]$ 
using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

show ?case proof(cases lnever isIntO ltr1)
  case True note ln3 = True note current = current True
  hence ln4: lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
  note ln34 = True this
  have ltrr1: ltrr1 = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrv1 ltrv1-lnever[OF current] by simp

have cltrv1: Van.lcompletedFrom sv1 (lltrv1 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-lltrv1[OF unW  $\Delta$  r ltr1 ln3 ltr2 ln4] .

show ?thesis apply(rule TwoFuncPred.sameFM-selectlmap-lfilter)
  unfolding ltrr1 apply simp
  using lS-lltrv1-ltr1[OF unW  $\Delta$  r ltr1 ln3 ltr2 ln4, of L]
  unfolding Van.lS[OF cltrv1] Opt.lS[OF ltr1(2)) .
next
  case False note ln3 = False
  hence ln4:  $\neg$  lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

  have ltr1  $\neq$   $[[s1]]$  using ln3 ltr1
  using Opt.final-not-isInt by auto
  hence llength ltr1 > Suc 0
  by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
  hence  $\neg$  finalO s1
  by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
  hence nf12:  $\neg$  finalV sv1  $\wedge$   $\neg$  finalV sv2
  using  $\Delta$  r(1) r(2) r(3) r(4) unW unwindCond-def by force

obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
  where  $\varphi\varphi$ :  $\varphi\varphi$  s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
  and  $\varphi'$ :  $\varphi'$   $\Delta$  w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO

```

```

and ltrr1: ltrr1 =
  lappend (llist-of trv1) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

using ltrv1-ltrv2-not-lnever[OF unw Δ r stat ltr1 ltr2 ln3 A34]
unfolding ltrr1 by blast
define ltrr1' where ltrr1': ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
unfolding ltrr1 ltrr1' ..
have ne1: trv1 ≠ [] ∨ w1' < w1
using  $\varphi'$  unfolding  $\varphi'$ -def ltrr1 by simp

show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of - w1])

apply(rule exI[of - tr1 ## s1']) apply(rule exI[of - w1']) apply(rule exI[of
- w1])
apply(rule exI[of - ltrr1']) apply(rule exI[of - s1'' $ ltr1'])
apply(intro conjI)
  subgoal unfolding nL .. subgoal unfolding nR ..
  subgoal using ltrr1 .
  subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
  subgoal by fact
  subgoal by simp
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def unfolding Van.S.map-filter Opt.S.map-filter
by simp
  subgoal apply(rule disjI1)
    apply(rule exI[of - w1']) apply(rule exI[of - w2'])
    apply(rule exI[of - s1''])
    apply(rule exI[of - s2'']) apply(rule exI[of - s2'' $ ltr2'])
    apply(rule exI[of - statAA])
    apply(rule exI[of - sv1'']) apply(rule exI[of - sv2''])
    apply(rule exI[of - statOO])
    apply(intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrr1' ..
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp

```

```

    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
  qed
  qed
}
thus ?thesis apply-
unfolding Van.lS[OF cltrv1] Opt.lS[OF ltr1(2)] apply(rule TwoFuncPred.sameFM-lmap-lfilter)
using assms by blast
qed

```

lemma lS-ltrv2-ltr2:

```

assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lS (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) = Opt.lS ltr2
proof-
have cltrv2: Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-ltrv2[OF assms] .
{fix nL nR ltrr2 ltr2
assume  $\exists w1 w2 s1 s2 ltr1 statA sv1 sv2 statO$ .
  nL = w2  $\wedge$  nR = w2  $\wedge$ 
  ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
   $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$ 
  reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
  (statA = Diff  $\longrightarrow$  statO = Diff)  $\wedge$ 
  Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$ 
  Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$ 
  Opt.lA ltr1 = Opt.lA ltr2
hence TwoFuncPred.sameFM isSecV isSecO getSecV getSecO nL nR ltrr2 ltr2
proof(coinduct rule: TwoFuncPred.sameFM.coinduct[of  $\lambda nL nR ltrr2 ltr2$ .
   $\exists w1 w2 s1 s2 ltr1 statA sv1 sv2 statO$ .
  nL = w2  $\wedge$  nR = w2  $\wedge$ 
  ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
   $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$ 
  reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
  (statA = Diff  $\longrightarrow$  statO = Diff)  $\wedge$ 
  Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$ 
  Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$ 
  Opt.lA ltr1 = Opt.lA ltr2])
case (2 nL nR ltrr2 ltr2)
then obtain w1 w2 s1 s2 ltr1 statA sv1 sv2 statO
where nL: nL = w2 and nR: nR = w2

```

```

and ltrr2: ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff → statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
by auto

have current: ltr1 ≠ [] ltr2 ≠ []
using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

show ?case proof(cases lnever isIntO ltr1)
  case True note ln3 = True note current = current True
  hence ln4: lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
  note ln34 = True this
  have ltrr2: ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr2 ltrv2-lnever[OF current] by simp

have cltrv2: Van.lcompletedFrom sv2 (lltrv2 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-lltrv2[OF unW Δ r ltr1 ln3 ltr2 ln4] .

show ?thesis apply(rule TwoFuncPred.sameFM-selectlmap-lfilter)
unfolding ltrr2 apply simp
using lS-lltrv2-ltr2[OF unW Δ r ltr1 ln3 ltr2 ln4, of L]
unfolding Van.lS[OF cltrv2] Opt.lS[OF ltr2(2)] .
next
case False note ln3 = False
hence ln4: ¬ lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

have ltr2 ≠ [[s2]] using ln4 ltr2
using Opt.final-not-isInt by auto
hence llength ltr2 > Suc 0
by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(2) enat-0-iff(2)
linorder-not-less llength-LNil llist-eq-cong ltr2(1) ltr2(2) nle-le not-less-zero)
hence ¬ finalO s2
by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(2) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr2(1))
hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
using Δ r(1) r(2) r(3) r(4) unW unwindCond-def by force

obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO

```

```

where  $\varphi\varphi$ :  $\varphi\varphi$   $s1$   $ltr1$   $s2$   $ltr2$   $tr1$   $s1'$   $s1''$   $ltr1'$   $tr2$   $s2'$   $s2''$   $ltr2'$ 
and  $\varphi'$ :  $\varphi' \Delta$   $w1$   $w2$   $w1'$   $w2'$   $statA$   $s1$   $tr1$   $s1'$   $s1''$   $s2$   $tr2$   $s2'$   $s2''$   $statAA$ 
 $statO$   $sv1$   $trv1$   $sv1''$   $sv2$   $trv2$   $sv2''$   $statOO$ 
and  $ltrr2$ :  $ltrr2 =$ 
 $lappend$  ( $l$ list-of  $trv2$ ) ( $ltrv2$  ( $w1'$ ,  $w2'$ ,  $s1''$ ,  $s1''$  $  $ltr1'$ ,  $s2''$ ,  $s2''$  $  $ltr2'$ ,  $statAA$ ,  $sv1''$ ,  $sv2''$ ,  $statOO$ ))

using  $ltrv1$ - $ltrv2$ -not-lnever[ $OF$   $unw$   $\Delta$   $r$   $stat$   $ltr1$   $ltr2$   $ln3$   $A34$ ]
unfolding  $ltrr2$  by  $blast$ 
define  $ltrr2'$  where  $ltrr2'$ :  $ltrr2' = ltrv2$  ( $w1'$ ,  $w2'$ ,  $s1''$ ,  $s1''$  $  $ltr1'$ ,  $s2''$ ,  $s2''$ 
$  $ltr2'$ ,  $statAA$ ,  $sv1''$ ,  $sv2''$ ,  $statOO$ )
have  $ltrr1$ :  $ltrr2 = lappend$  ( $l$ list-of  $trv2$ )  $ltrr2'$ 
unfolding  $ltrr2$   $ltrr2'$  ..
have  $ne2$ :  $trv2 \neq [] \vee w2' < w2$ 
using  $\varphi'$  unfolding  $\varphi'$ -def  $ltrr2$  by  $simp$ 

show ?thesis
apply(rule  $TwoFuncPred$ .sameFM-selectlappend)
apply(rule  $exI$ [of -  $trv2$ ]) apply(rule  $exI$ [of -  $w2'$ ]) apply(rule  $exI$ [of -  $w2$ ])

apply(rule  $exI$ [of -  $tr2$  ##  $s2'$ ]) apply(rule  $exI$ [of -  $w2'$ ]) apply(rule  $exI$ [of
-  $w2$ ])
apply(rule  $exI$ [of -  $ltrr2'$ ]) apply(rule  $exI$ [of -  $s2''$  $  $ltr2'$ ])
apply(intro  $conjI$ )
subgoal unfolding  $nL$  .. subgoal unfolding  $nR$  ..
subgoal using  $ltrr1$  .
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by  $simp$ 
subgoal by  $fact$ 
subgoal by  $simp$ 
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def unfolding  $Van.S$ .map-filter  $Opt.S$ .map-filter
by  $simp$ 
subgoal apply(rule  $disjI1$ )
apply(rule  $exI$ [of -  $w1'$ ]) apply(rule  $exI$ [of -  $w2'$ ])
apply(rule  $exI$ [of -  $s1''$ ])
apply(rule  $exI$ [of -  $s2''$ ]) apply(rule  $exI$ [of -  $s1''$  $  $ltr1'$ ])
apply(rule  $exI$ [of -  $statAA$ ])
apply(rule  $exI$ [of -  $sv1''$ ]) apply(rule  $exI$ [of -  $sv2''$ ])
apply(rule  $exI$ [of -  $statOO$ ])
apply(intro  $conjI$ )
subgoal .. subgoal ..
subgoal unfolding  $ltrr2'$  ..
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by  $simp$ 
subgoal using  $r(1)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by ( $metis$   $Opt$ .reach-validFromS-reach  $append$ -is- $Nil$ -conv  $last$ -snoc
not-Cons-self2)
subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
by ( $metis$   $Opt$ .reach-validFromS-reach  $append$ -is- $Nil$ -conv  $last$ -snoc
not-Cons-self2)
subgoal using  $r(3)$   $\varphi'$  unfolding  $\varphi'$ -def
by ( $metis$   $Van$ .reach-validFromS-reach  $snoc$ -eq-iff-butlast)

```



```

    subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
    by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
  qed
  qed
}
thus ?thesis apply-
unfolding Van.lS[OF cltrv2] Opt.lS[OF ltr2(2)] apply(rule TwoFuncPred.sameFM-lmap-lfilter)
using assms by blast
qed

```

lemma *lA-ltrv1-ltrv2*:

```

assumes unw: unwindCond  $\Delta$ 
and  $\Delta$ :  $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$ 
and r: reachO s1 reachO s2 reachV sv1 reachV sv2
and stat: statA = Diff  $\longrightarrow$  statO = Diff
and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
and A34: Opt.lA ltr1 = Opt.lA ltr2
shows Van.lA (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) =
  Van.lA (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
proof-
have cltrv1: Van.lcompletedFrom sv1 (ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-ltrv1[OF assms] .
have cltrv2: Van.lcompletedFrom sv2 (ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-ltrv2[OF assms] .
{fix nL nR ltrr1 ltrr2
assume  $\exists w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO.$ 
  nL = w1  $\wedge$  nR = w2  $\wedge$ 
  ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
  ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 
   $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO \wedge$ 
  reachO s1  $\wedge$  reachO s2  $\wedge$  reachV sv1  $\wedge$  reachV sv2  $\wedge$ 
  (statA = Diff  $\longrightarrow$  statO = Diff)  $\wedge$ 
  Opt.lvalidFromS s1 ltr1  $\wedge$  Opt.lcompletedFrom s1 ltr1  $\wedge$ 
  Opt.lvalidFromS s2 ltr2  $\wedge$  Opt.lcompletedFrom s2 ltr2  $\wedge$ 
  Opt.lA ltr1 = Opt.lA ltr2
hence TwoFuncPred.sameFM isIntV isIntV getActV getActV nL nR ltrr1 ltrr2
proof(coinduct rule: TwoFuncPred.sameFM.coinduct[of  $\lambda nL nR ltrr1 ltrr2.$ 
   $\exists w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO.$ 
  nL = w1  $\wedge$  nR = w2  $\wedge$ 
  ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)  $\wedge$ 

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    ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) ∧
    Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO ∧
    reachO s1 ∧ reachO s2 ∧ reachV sv1 ∧ reachV sv2 ∧
    (statA = Diff → statO = Diff) ∧
    Opt.lvalidFromS s1 ltr1 ∧ Opt.lcompletedFrom s1 ltr1 ∧
    Opt.lvalidFromS s2 ltr2 ∧ Opt.lcompletedFrom s2 ltr2 ∧
    Opt.lA ltr1 = Opt.lA ltr2])
  case (2 nL nR ltrr1 ltrr2)
  then obtain w1 w2 s1 ltr1 s2 ltr2 statA sv1 sv2 statO
  where nL: nL = w1 and nR: nR = w2
  and ltrr1: ltrr1 = ltrv1 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and ltrr2: ltrr2 = ltrv2 (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)
  and Δ: Δ ∞ w1 w2 s1 s2 statA sv1 sv2 statO
  and r: reachO s1 reachO s2 reachV sv1 reachV sv2
  and stat: statA = Diff → statO = Diff
  and ltr1: Opt.lvalidFromS s1 ltr1 Opt.lcompletedFrom s1 ltr1
  and ltr2: Opt.lvalidFromS s2 ltr2 Opt.lcompletedFrom s2 ltr2
  and A34: Opt.lA ltr1 = Opt.lA ltr2
  by auto

  have current: ltr1 ≠ [] ltr2 ≠ []
  using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

  show ?case proof (cases lnever isIntO ltr1)
  case True note ln3 = True note current = current True
  hence ln4: lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
    ltr2(2))
  note ln34 = True this
  have ltrr1: ltrr1 = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
    statO)
  unfolding ltrr1 ltrv1-lnever[OF current] by simp
  have ltrr2: ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
    statO)
  unfolding ltrr2 ltrv2-lnever[OF current] by simp

  have cltrv1: Van.lcompletedFrom sv1 (lltrv1 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-lltrv1[OF unW Δ r ltr1 ln3 ltr2 ln4] .
  have cltrv2: Van.lcompletedFrom sv2 (lltrv2 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
  using lcompletedFrom-lltrv2[OF unW Δ r ltr1 ln3 ltr2 ln4] .

  show ?thesis apply (rule TwoFuncPred.sameFM-selectlmap-lfilter)
  unfolding ltrr1 ltrr2 apply simp
  using lA-lltrv1-lltrv2[OF unW Δ r ltr1 ln3 ltr2 ln4, of L]
  unfolding Van.lA[OF cltrv1] Van.lA[OF cltrv2] .
  next
  case False note ln3 = False
  hence ln4: ¬ lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34

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ltr2(2))

```
have ltr1 ≠ [[s1]] using ln3 ltr1
using Opt.final-not-isInt by auto
hence llength ltr1 > Suc 0
by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
hence ¬ finalO s1
by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
using Δ r(1) r(2) r(3) r(4) unW unwindCond-def by force

obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
where φφ: φφ s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'
and φ': φ' Δ w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA
statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO
and ltrr1: ltrr1 =
lappend (llist-of trv1) (ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

and ltrr2: ltrr2 =
lappend (llist-of trv2) (ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2'' $ ltr2',statAA,sv1'',sv2'',statOO))

using ltrv1-ltrv2-not-lnever[OF unW Δ r stat ltr1 ltr2 ln3 A34]
unfolding ltrr1 ltrr2 by blast
define ltrr1' where ltrr1': ltrr1' = ltrv1 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
have ltrr1: ltrr1 = lappend (llist-of trv1) ltrr1'
unfolding ltrr1 ltrr1' ..
have ne1: trv1 ≠ [] ∨ w1' < w1
using φ' unfolding φ'-def ltrr1 by simp
define ltrr2' where ltrr2': ltrr2' = ltrv2 (w1',w2',s1'',s1'' $ ltr1',s2'',s2''
$ ltr2',statAA,sv1'',sv2'',statOO)
have ltrr2: ltrr2 = lappend (llist-of trv2) ltrr2'
unfolding ltrr2 ltrr2' ..
have ne2: trv2 ≠ [] ∨ w2' < w2
using φ' unfolding φ'-def ltrr1 by simp

show ?thesis
apply(rule TwoFuncPred.sameFM-selectlappend)
apply(rule exI[of - trv1]) apply(rule exI[of - w1']) apply(rule exI[of - w1])

apply(rule exI[of - trv2]) apply(rule exI[of - w2']) apply(rule exI[of - w2])

apply(rule exI[of - ltrr1']) apply(rule exI[of - ltrr2'])
apply(intro conjI)
subgoal unfolding nL .. subgoal unfolding nR ..
```

```

subgoal using ltrr1 .
subgoal using ltrr2 .
subgoal by fact subgoal by fact
subgoal using  $\varphi'$  unfolding  $\varphi'$ -def unfolding Van.A.map-filter by simp
subgoal apply (rule disjI1)
  apply (rule exI[of - w1']) apply (rule exI[of - w2'])
  apply (rule exI[of - s1'']) apply (rule exI[of - s1'' $ ltr1'])
  apply (rule exI[of - s2'']) apply (rule exI[of - s2'' $ ltr2'])
  apply (rule exI[of - statAA])
  apply (rule exI[of - sv1'']) apply (rule exI[of - sv2''])
  apply (rule exI[of - statOO])
  apply (intro conjI)
    subgoal .. subgoal ..
    subgoal unfolding ltrr1' ..
    subgoal unfolding ltrr2' ..
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using r(1)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def
      by (metis Opt.reach-validFromS-reach append-is-Nil-conv last-snoc
not-Cons-self2)
    subgoal using r(3)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using r(4)  $\varphi'$  unfolding  $\varphi'$ -def
      by (metis Van.reach-validFromS-reach snoc-eq-iff-butlast)
    subgoal using  $\varphi'$  unfolding  $\varphi'$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp
    subgoal using r(2)  $\varphi\varphi$  unfolding  $\varphi\varphi$ -def by simp . .
  qed
qed
}
thus ?thesis apply-
unfolding Van.lA[OF cltrv1] Van.lA[OF cltrv2] apply (rule TwoFuncPred.sameFM-lmap-lfilter)
using assms by blast
qed

```

lemma *lO-ltrv1-ltrv2*:
assumes *unw*: *unwindCond* Δ
and Δ : $\Delta \infty w1 w2 s1 s2 statA sv1 sv2 statO$
and *r*: *reachO* *s1* *reachO* *s2* *reachV* *sv1* *reachV* *sv2*
and *stat*: *statA* = *Diff* \longrightarrow *statO* = *Diff*
and *ltr1*: *Opt.lvalidFromS* *s1* *ltr1* *Opt.lcompletedFrom* *s1* *ltr1*
and *ltr2*: *Opt.lvalidFromS* *s2* *ltr2* *Opt.lcompletedFrom* *s2* *ltr2*

and $A34$: $Opt.lA\ ltr1 = Opt.lA\ ltr2$
and $O12$: $Van.lO\ (ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)) =$
 $Van.lO\ (ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))$
and stO : $statO = Eq$
shows $Opt.lO\ ltr1 = Opt.lO\ ltr2$
proof –
have $cltrv1$: $Van.lcompletedFrom\ sv1\ (ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))$
using $lcompletedFrom-ltrv1[OF\ assms(1-12)]$.
have $cltrv2$: $Van.lcompletedFrom\ sv2\ (ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))$
using $lcompletedFrom-ltrv2[OF\ assms(1-12)]$.
{fix $nL\ nR\ ltr1\ ltr2$
assume $\exists\ ltrr1\ ltrr2\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$.
 $ltrr1 = ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$
 $ltrr2 = ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$
 $\Delta \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge$
 $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge$
 $Opt.lA\ ltr1 = Opt.lA\ ltr2 \wedge$
 $Van.lO\ ltrr1 = Van.lO\ ltrr2 \wedge$
 $statO = Eq$
hence $TwoFuncPred.sameFM\ isIntO\ isIntO\ getObsO\ getObsO\ nL\ nR\ ltr1\ ltr2$
proof(*coinduct rule: TwoFuncPred.sameFM.coinduct[of $\lambda nL\ nR\ ltr1\ ltr2$.*
 $\exists\ ltrr1\ ltrr2\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$.
 $ltrr1 = ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$
 $ltrr2 = ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO) \wedge$
 $\Delta \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \wedge$
 $reachO\ s1 \wedge reachO\ s2 \wedge reachV\ sv1 \wedge reachV\ sv2 \wedge$
 $(statA = Diff \longrightarrow statO = Diff) \wedge$
 $Opt.lvalidFromS\ s1\ ltr1 \wedge Opt.lcompletedFrom\ s1\ ltr1 \wedge$
 $Opt.lvalidFromS\ s2\ ltr2 \wedge Opt.lcompletedFrom\ s2\ ltr2 \wedge$
 $Opt.lA\ ltr1 = Opt.lA\ ltr2 \wedge$
 $Van.lO\ ltrr1 = Van.lO\ ltrr2 \wedge$
 $statO = Eq$)
case ($2\ nL\ nR\ ltr1\ ltr2$)
then obtain $ltrr1\ ltrr2\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ **where**
 $ltrr11$: $ltrr1 = ltrv1\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)$
and $ltrr22$: $ltrr2 = ltrv2\ (w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO)$
and Δ : $\Delta \infty\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$
and r : $reachO\ s1\ reachO\ s2\ reachV\ sv1\ reachV\ sv2$
and $stat$: $statA = Diff \longrightarrow statO = Diff$
and $ltr1$: $Opt.lvalidFromS\ s1\ ltr1\ Opt.lcompletedFrom\ s1\ ltr1$
and $ltr2$: $Opt.lvalidFromS\ s2\ ltr2\ Opt.lcompletedFrom\ s2\ ltr2$
and $A34$: $Opt.lA\ ltr1 = Opt.lA\ ltr2$
and $O12$: $Van.lO\ ltrr1 = Van.lO\ ltrr2$
and stO : $statO = Eq$
by *auto*

```

have stA: statA = Eq using stat stO
using status.exhaust by blast

have current: ltr1 ≠ [] ltr2 ≠ []
using ltr1(2) ltr2(2) unfolding Opt.lcompletedFrom-def by auto

show ?case proof(cases lnever isIntO ltr1)
  case True note ln3 = True note current = current True
  hence ln4: lnever isIntO ltr2
  by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))
  note ln34 = True this
  have ltrr1: ltrr1 = lltrv1 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr11 ltrv1-lnever[OF current] by simp
  have ltrr2: ltrr2 = lltrv2 (L, w1, w2, s1, ltr1, s2, ltr2, statA, sv1, sv2,
statO)
  unfolding ltrr22 ltrv2-lnever[OF current] by simp

have clltrv1: Van.lcompletedFrom sv1 (lltrv1 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-lltrv1[OF unW Δ r ltr1 ln3 ltr2 ln4] .
have clltrv2: Van.lcompletedFrom sv2 (lltrv2 (L,w1,w2,s1,ltr1,s2,ltr2,statA,sv1,sv2,statO))
using lcompletedFrom-lltrv2[OF unW Δ r ltr1 ln3 ltr2 ln4] .

show ?thesis apply(rule TwoFuncPred.sameFM-selectlmap-lfilter)
unfolding Opt.lO[OF ltr1(2)] Opt.lO[OF ltr2(2)]
by (metis ln3 ln4 lnever-LNil-lfilter')

next
case False note ln3 = False
hence ln4: ¬ lnever isIntO ltr2
by (metis Opt.lA lfiltermap-LNil-never lfiltermap-lmap-lfilter ltr1(2) A34
ltr2(2))

have ltr1 ≠ [s1] using ln3 ltr1
using Opt.final-not-isInt by auto
hence llength ltr1 > Suc 0
by (metis Opt.lcompletedFrom-singl Suc-ile-eq current(1) enat-0-iff(1)
linorder-not-less llength-LNil llist-eq-cong ltr1(1) ltr1(2) nle-le not-less-zero)
hence ¬ finalO s1
by (metis Opt.final-def Opt.lvalidFromS-Cons-iff current(1) eSuc-enat enat-0

linorder-neq-iff llength-LCons llength-LNil llist.exhaust-sel ltr1(1))
hence nf12: ¬ finalV sv1 ∧ ¬ finalV sv2
using Δ r(1) r(2) r(3) r(4) unW unwindCond-def by force

obtain w1' w2' tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2' trv1 sv1'' trv2 sv2'' statAA
statOO
where φφ: φφ s1 ltr1 s2 ltr2 tr1 s1' s1'' ltr1' tr2 s2' s2'' ltr2'

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and  $\varphi'$ :  $\varphi' \Delta w1 w2 w1' w2' statA s1 tr1 s1' s1'' s2 tr2 s2' s2'' statAA$ 
 $statO sv1 trv1 sv1'' sv2 trv2 sv2'' statOO$ 
and  $ltrr1$ :  $ltrr1 =$ 
 $lappend (l\text{list-of } trv1) (ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$ 

and  $ltrr2$ :  $ltrr2 =$ 
 $lappend (l\text{list-of } trv2) (ltrv2 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2'' \$ ltr2', statAA, sv1'', sv2'', statOO))$ 

using  $ltrv1$ - $ltrv2$ - $\text{not-lnever}$ [ $OF unw \Delta r stat ltr1 ltr2 ln3 A34$ ]
unfolding  $ltrr11$   $ltrr22$  by  $blast$ 
define  $ltrr1'$  where  $ltrr1'$ :  $ltrr1' = ltrv1 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2''$ 
 $\$ ltr2', statAA, sv1'', sv2'', statOO)$ 
have  $ltrr1$ :  $ltrr1 = lappend (l\text{list-of } trv1) ltrr1'$ 
unfolding  $ltrr1$   $ltrr1'$  ..
have  $ne1$ :  $trv1 \neq [] \vee w1' < w1$ 
using  $\varphi'$  unfolding  $\varphi'$ - $\text{def } ltrr1$  by  $simp$ 
define  $ltrr2'$  where  $ltrr2'$ :  $ltrr2' = ltrv2 (w1', w2', s1'', s1'' \$ ltr1', s2'', s2''$ 
 $\$ ltr2', statAA, sv1'', sv2'', statOO)$ 
have  $ltrr2$ :  $ltrr2 = lappend (l\text{list-of } trv2) ltrr2'$ 
unfolding  $ltrr2$   $ltrr2'$  ..
have  $ne2$ :  $trv2 \neq [] \vee w2' < w2$ 
using  $\varphi'$  unfolding  $\varphi'$ - $\text{def } ltrr1$  by  $simp$ 

have  $ltr1$ - $eq$ :  $ltr1 = lappend (l\text{list-of } (tr1 \#\# s1')) (s1'' \$ ltr1')$ 
and  $ltr2$ - $eq$ :  $ltr2 = lappend (l\text{list-of } (tr2 \#\# s2')) (s2'' \$ ltr2')$  using  $\varphi\varphi$ 
unfolding  $\varphi\varphi$ - $\text{def by auto}$ 

have  $sst$ :  $statOO = Diff \longleftrightarrow Van.O (trv1 \#\# sv1'') \neq Van.O (trv2 \#\#$ 
 $sv2'')$ 
 $statA = Eq \implies statAA = Diff \longleftrightarrow Opt.O ((tr1 \#\# s1') \#\# s1'') \neq Opt.O$ 
 $((tr2 \#\# s2') \#\# s2'')$ 
 $statO = Diff \implies statOO = Diff$ 
 $statAA = Diff \implies statOO = Diff$ 
using  $\varphi'$   $stO$  unfolding  $\varphi'$ - $\text{def by auto}$ 

have  $Atrv12'$ :  $Van.A (trv1 \#\# sv1'') = Van.A (trv2 \#\# sv2'')$ 
using  $\varphi'$  unfolding  $\varphi'$ - $\text{def by auto}$ 

have  $\Delta'$ :  $\Delta \infty w1' w2' s1'' s2'' statAA sv1'' sv2'' statOO$ 
using  $\varphi'$  unfolding  $\varphi'$ - $\text{def by auto}$ 

have  $ultrv1$ :  $Van.l\text{validFromS } sv1$   $ltrr1$ 
unfolding  $ltrr11$  using  $l\text{validFromS-ltrv1}$ 
using  $A34 \Delta ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw$ 
by  $blast$ 
have  $cltrv1$ :  $Van.l\text{completedFrom } sv1$   $ltrr1$ 
unfolding  $ltrr11$  using  $l\text{completedFrom-ltrv1}$ 
using  $A34 \Delta ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw$ 
by  $blast$ 

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```

have ultrv2: Van.lvalidFromS sv2 ltrr2
unfolding ltrr22 using lvalidFromS-ltrv2
using A34  $\Delta$  ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw
by blast
have cltrv2: Van.lcompletedFrom sv2 ltrr2
unfolding ltrr22 using lcompletedFrom-ltrv2
using A34  $\Delta$  ltr1(1) ltr1(2) ltr2(1) ltr2(2) r(1) r(2) r(3) r(4) stat unw
by blast

have Oltrr1: Van.lO ltrr1 = lmap getObsV (lfilter isIntV ltrr1)
using Van.lO[OF cltrv1] .
have Oltrr2: Van.lO ltrr2 = lmap getObsV (lfilter isIntV ltrr2)
using Van.lO[OF cltrv2] .

have cltrv1': Van.lcompletedFrom (lastt sv1 trv1) ltrr1'
using cltrv1 unfolding ltrr1 Van.lcompletedFrom-def Van.final-def apply
simp
using  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.validS-validTrans Van.validFromS-def lappend-LNil2 last-appendR
lfinite-llist-of list-of-lappend
list-of-llist-of llist-of.simps(1) snoc-eq-iff-butlast)

have cltrv2': Van.lcompletedFrom (lastt sv2 trv2) ltrr2'
using cltrv2 unfolding ltrr2 Van.lcompletedFrom-def Van.final-def apply
simp
using  $\varphi'$  unfolding  $\varphi'$ -def
by (metis Van.validS-validTrans Van.validFromS-def lappend-LNil2 last-appendR
lfinite-llist-of list-of-lappend
list-of-llist-of llist-of.simps(1) snoc-eq-iff-butlast)

have Oltrr1': Van.lO ltrr1' = lmap getObsV (lfilter isIntV ltrr1')
using Van.lO[OF cltrv1'] .
have Oltrr2': Van.lO ltrr2' = lmap getObsV (lfilter isIntV ltrr2')
using Van.lO[OF cltrv2'] .

have Van.O (trv1 ## sv1'') = Van.O (trv2 ## sv2'')  $\wedge$ 
Van.lO ltrr1' = Van.lO ltrr2'
using Atrv12' O12 Oltrr1' Oltrr2' unfolding Oltrr1 Oltrr2 unfolding ltrr1
ltrr2
unfolding Van.A.map-filter Van.O.map-filter Van.lO.lmap-lfilter apply simp

unfolding lmap-lappend-distrib apply simp apply(subst (asm) lappend-llist-of-inj)

using map-eq-imp-length-eq by auto
hence O12'': Van.O (trv1 ## sv1'') = Van.O (trv2 ## sv2'')
and O12': Van.lO ltrr1' = Van.lO ltrr2' by auto

have stOO: statOO = Eq using O12'' sst by(cases statOO, auto)

```



```

have  $O34'$ :  $Opt.O ((tr1 \#\# s1') \#\# s1'') = Opt.O ((tr2 \#\# s2') \#\#$ 
 $s2'')$ 
using  $stOO$   $sst(2)$   $sst(4)$   $stA$  by  $blast$ 

hence  $s14'$ :  $getObsO s1' = getObsO s2'$ 
using  $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$   $Opt.O.map-filter$  by ( $simp$   $add$ :  $never-Nil-filter$ )

show  $?thesis$ 
apply( $rule$   $TwoFuncPred.sameFM-selectlappend$ )
apply( $rule$   $exI[of - tr1 \#\# s1']$ ) apply( $rule$   $exI[of - undefined]$ ) apply( $rule$ 
 $exI[of - nL]$ )
apply( $rule$   $exI[of - tr2 \#\# s2']$ ) apply( $rule$   $exI[of - undefined]$ ) apply( $rule$ 
 $exI[of - nR]$ )
apply( $rule$   $exI[of - s1'' \$ ltr1']$ ) apply( $rule$   $exI[of - s2'' \$ ltr2']$ )
apply( $intro$   $conjI$ )
subgoal .. subgoal ..
subgoal by  $fact$  subgoal by  $fact$ 
subgoal by  $simp$  subgoal by  $simp$ 
subgoal using  $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$  unfolding  $Opt.O.map-filter$ 
by ( $simp$   $add$ :  $s14'$   $never-Nil-filter$ )
subgoal apply( $rule$   $disjI1$ )
apply( $rule$   $exI[of - ltrr1']$ ) apply( $rule$   $exI[of - ltrr2']$ )
apply( $rule$   $exI[of - w1']$ ) apply( $rule$   $exI[of - w2']$ )
apply( $rule$   $exI[of - s1'']$ ) apply( $rule$   $exI[of - s2'']$ )
apply( $rule$   $exI[of - statAA]$ )
apply( $rule$   $exI[of - sv1'']$ ) apply( $rule$   $exI[of - sv2'']$ )
apply( $rule$   $exI[of - statOO]$ )
apply( $intro$   $conjI$ )
subgoal unfolding  $ltrr1'$  ..
subgoal unfolding  $ltrr2'$  ..
subgoal using  $\varphi'$  unfolding  $\varphi'$ - $def$  by  $simp$ 
subgoal using  $r(1)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$ 
by ( $metis$   $Opt.reach-validFromS-reach$   $append-is-Nil-conv$   $last-snoc$ 
 $not-Cons-self2$ )
subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$ 
by ( $metis$   $Opt.reach-validFromS-reach$   $append-is-Nil-conv$   $last-snoc$ 
 $not-Cons-self2$ )
subgoal using  $r(3)$   $\varphi'$  unfolding  $\varphi'$ - $def$ 
by ( $metis$   $Van.reach-validFromS-reach$   $snoc-eq-iff-butlast$ )
subgoal using  $r(4)$   $\varphi'$  unfolding  $\varphi'$ - $def$ 
by ( $metis$   $Van.reach-validFromS-reach$   $snoc-eq-iff-butlast$ )
subgoal using  $\varphi'$  unfolding  $\varphi'$ - $def$  by  $simp$ 
subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$  by  $simp$ 
subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$  by  $simp$ 
subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$  by  $simp$ 
subgoal using  $r(2)$   $\varphi\varphi$  unfolding  $\varphi\varphi$ - $def$  by  $simp$ 
subgoal by  $fact$ 

```

```

      subgoal by fact . .
    qed
  qed
}
thus ?thesis
unfolding Opt.lO[OF ltr1(2)] Opt.lO[OF ltr2(2)] apply(rule TwoFuncPred.sameFM-lmap-lfilter[where
wL = undefined and wR = undefined])
using assms by blast
qed
end

```

```

theorem unwind-lrsecure:
assumes init: initCond  $\Delta$  and unwind: unwindCond  $\Delta$ 
shows lrsecure
unfolding lrsecure-def2 proof clarify
  fix s1 tr1 s2 tr2
  assume 3: istateO s1 Opt.lvalidFromS s1 tr1 lcompletedFromO s1 tr1
  and 4: istateO s2 Opt.lvalidFromS s2 tr2 lcompletedFromO s2 tr2
  and A34: Opt.lA tr1 = Opt.lA tr2 and O34: Opt.lO tr1  $\neq$  Opt.lO tr2
  obtain sv1 sv2 where
  isv12: istateV sv1 istateV sv2 and c12: corrState sv1 s1 corrState sv2 s2
  and  $\Delta$ :  $\Delta \infty \infty \infty s1 s2 Eq sv1 sv2 Eq$ 
  using init 3 4 unfolding initCond-def by blast
  have r: reachV sv1 reachV sv2 reachO s1 reachO s2
  by (auto simp: Van.Istate isv12 Opt.Istate 3 4)
  note all = 3 4 A34 isv12  $\Delta$  unwind r
  show  $\exists sv1 trv1 sv2 trv2$ .
    istateV sv1  $\wedge$  istateV sv2  $\wedge$  corrState sv1 s1  $\wedge$  corrState sv2 s2  $\wedge$ 
    Van.lvalidFromS sv1 trv1  $\wedge$  lcompletedFromV sv1 trv1  $\wedge$  Van.lvalidFromS sv2
  trv2  $\wedge$ 
    lcompletedFromV sv2 trv2  $\wedge$  Van.lS trv1 = Opt.lS tr1  $\wedge$  Van.lS trv2 = Opt.lS
  tr2  $\wedge$ 
    Van.lA trv1 = Van.lA trv2  $\wedge$  Van.lO trv1  $\neq$  Van.lO trv2
  apply(rule exI[of - sv1])
  apply(rule exI[of - ltrv1  $\Delta (\infty, \infty, s1, tr1, s2, tr2, Eq, sv1, sv2, Eq)$ ])
  apply(rule exI[of - sv2])
  apply(rule exI[of - ltrv2  $\Delta (\infty, \infty, s1, tr1, s2, tr2, Eq, sv1, sv2, Eq)$ ])
  apply(intro conjI)
  subgoal by fact subgoal by fact subgoal by fact subgoal by fact
  subgoal apply(rule lvalidFromS-ltrv1) using all by auto
  subgoal apply(rule lcompletedFrom-ltrv1) using all by auto
  subgoal apply(rule lvalidFromS-ltrv2) using all by auto
  subgoal apply(rule lcompletedFrom-ltrv2) using all by auto
  subgoal apply(rule lS-ltrv1-ltr1) using all by auto
  subgoal apply(rule lS-ltrv2-ltr2) using all by auto

```

subgoal apply(rule lA-ltrv1-ltrv2) **using all by auto**
subgoal using O34 **apply– apply**(erule contrapos-nn)
apply(rule lO-ltrv1-ltrv2) **using all by auto** .
qed

4.4 Compositional unwinding

We allow networks of unwinding relations that unwind into each other, which offer a compositional alternative to monolithic unwinding.

definition *unwindIntoCond* ::

$(\text{enat} \Rightarrow \text{enat} \Rightarrow \text{enat} \Rightarrow \text{'stateO} \Rightarrow \text{'stateO} \Rightarrow \text{status} \Rightarrow \text{'stateV} \Rightarrow \text{'stateV} \Rightarrow \text{status} \Rightarrow \text{bool}) \Rightarrow$
 $(\text{enat} \Rightarrow \text{enat} \Rightarrow \text{enat} \Rightarrow \text{'stateO} \Rightarrow \text{'stateO} \Rightarrow \text{status} \Rightarrow \text{'stateV} \Rightarrow \text{'stateV} \Rightarrow \text{status} \Rightarrow \text{bool})$
 $\Rightarrow \text{bool}$

where

$\text{unwindIntoCond } \Delta \Delta' \equiv \forall w w1 w2 s1 s2 \text{statA sv1 sv2 statO.}$
 $\text{reachO } s1 \wedge \text{reachO } s2 \wedge \text{reachV } sv1 \wedge \text{reachV } sv2 \wedge$
 $\Delta w w1 w2 s1 s2 \text{statA sv1 sv2 statO} \longrightarrow$
 $(\text{finalO } s1 \longleftrightarrow \text{finalO } s2) \wedge (\text{finalV } sv1 \longleftrightarrow \text{finalO } s1) \wedge (\text{finalV } sv2 \longleftrightarrow \text{finalO } s2)$
 \wedge
 $(\text{statA} = \text{Eq} \longrightarrow (\text{isIntO } s1 \longleftrightarrow \text{isIntO } s2))$
 \wedge
 $((\exists v < w. \text{proact } \Delta' v w1 w2 s1 s2 \text{statA sv1 sv2 statO})$
 \vee
 $\text{react } \Delta' w1 w2 s1 s2 \text{statA sv1 sv2 statO})$

theorem *distrib-unwind-lrsecure*:

assumes $m: 0 < m$ **and** $\text{nxt}: \bigwedge i. i < (m::\text{nat}) \implies \text{nxt } i \subseteq \{0..<m\}$

and $\text{init}: \text{initCond } (\Delta s 0)$

and $\text{step}: \bigwedge i. i < m \implies$

$\text{unwindIntoCond } (\Delta s i) (\lambda w w1 w2 s1 s2 \text{statA sv1 sv2 statO.}$

$\exists j \in \text{nxt } i. \Delta s j w w1 w2 s1 s2 \text{statA sv1 sv2 statO})$

shows *lrsecure*

proof –

define Δ **where** $D: \Delta \equiv \lambda w w1 w2 s1 s2 \text{statA sv1 sv2 statO.} \exists i < m. \Delta s i w w1 w2 s1 s2 \text{statA sv1 sv2 statO}$

have $i: \text{initCond } \Delta$ **using** $\text{init } m$ **unfolding** *initCond-def* **by** *meson*

have $c: \text{unwindCond } \Delta$ **unfolding** *unwindCond-def* **apply**(*intro allI impI allI*)

apply(*subst (asm) D*) **apply** (*elim exE conjE*)

subgoal for $w w1 w2 s1 s2 \text{statA sv1 sv2 statO } i$

apply(*frule step*) **unfolding** *unwindIntoCond-def*

apply(erule *allE*[of - w]) **apply**(erule *allE*[of - $w1$]) **apply**(erule *allE*[of - $w2$])

apply(erule *allE*[of - $s1$]) **apply**(erule *allE*[of - $s2$]) **apply**(erule *allE*[of - statA])

apply(erule *allE*[of - sv1]) **apply**(erule *allE*[of - sv2]) **apply**(erule *allE*[of - statO])

```

apply simp apply(elim conjE)
apply(erule disjE)
  subgoal apply(rule disjI1)
  subgoal apply(elim exE conjE) subgoal for v
  apply(rule exI[of - v], simp)
  apply(rule proact-mono[of λw w1 w2 s1 s2 statA sv1 sv2 statO. ∃j∈next i.
Δs j w w1 w2 s1 s2 statA sv1 sv2 statO])
  subgoal unfolding le-fun-def D by simp (meson atLeastLessThan-iff next
subsetD)
  subgoal . . . .
  subgoal apply(rule disjI2)
  apply(rule match-mono[of λw w1 w2 s1 s2 statA sv1 sv2 statO. ∃j∈next i.
Δs j w w1 w2 s1 s2 statA sv1 sv2 statO])
  subgoal unfolding le-fun-def D by simp (meson atLeastLessThan-iff next
subsetD)
  subgoal . . . .
  show ?thesis using unwind-lrsecure[OF i c] .
qed

```

lemma *unwindIntoCond-simpleI*:

assumes

```

 $\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$ 
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$ 
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ 
 $\implies$ 
 $(finalO s1 \longleftrightarrow finalO s2) \wedge (finalV sv1 \longleftrightarrow finalO s1) \wedge (finalV sv2 \longleftrightarrow finalO$ 
 $s2)$ 

```

and

```

 $\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$ 
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$ 
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$ 
 $statA = Eq$ 
 $\implies$ 

```

```

 $isIntO s1 \longleftrightarrow isIntO s2$ 

```

```

 $\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$ 
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$ 
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ 
 $\implies$ 

```

```

 $react \Delta' w1 w2 s1 s2 statA sv1 sv2 statO$ 

```

shows *unwindIntoCond* $\Delta \Delta'$

using *assms* **unfolding** *unwindIntoCond-def* **by** *auto*

lemma *unwindIntoCond-simpleI2*:

assumes

```

 $\bigwedge w w1 w2 s1 s2 statA sv1 sv2 statO.$ 
 $reachO s1 \implies reachO s2 \implies reachV sv1 \implies reachV sv2 \implies$ 
 $\Delta w w1 w2 s1 s2 statA sv1 sv2 statO$ 
 $\implies$ 

```

$(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

and

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$

$statA = Eq$

\implies

$isIntO\ s1 \longleftrightarrow isIntO\ s2$

and

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

\implies

$(\exists v < w. proact\ \Delta'\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO)$

shows $unwindIntoCond\ \Delta\ \Delta'$

using *assms* **unfolding** *unwindIntoCond-def* **by** *auto*

lemma *unwindIntoCond-simpleIB*:

assumes

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

\implies

$(finalO\ s1 \longleftrightarrow finalO\ s2) \wedge (finalV\ sv1 \longleftrightarrow finalO\ s1) \wedge (finalV\ sv2 \longleftrightarrow finalO\ s2)$

and

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$

$statA = Eq$

\implies

$isIntO\ s1 \longleftrightarrow isIntO\ s2$

and

$\bigwedge w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$reachO\ s1 \implies reachO\ s2 \implies reachV\ sv1 \implies reachV\ sv2 \implies$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

\implies

$(\exists v < w. proact\ \Delta'\ v\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO) \vee react\ \Delta'\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

shows $unwindIntoCond\ \Delta\ \Delta'$

using *assms* **unfolding** *unwindIntoCond-def* **by** *auto*

definition *oor where*

$oor\ \Delta\ \Delta_2 \equiv \lambda w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO.$

$\Delta\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \vee \Delta_2\ w\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$

lemma *oorI1*:

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor \Delta \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor-def* by *simp*

lemma *oorI2*:

$\Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor \Delta \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor-def* by *simp*

definition *oor3* **where**

$oor3 \Delta \Delta_2 \Delta_3 \equiv \lambda w w1 w2 s1 s2 statA sv1 sv2 statO.$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO$

lemma *oor3I1*:

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor3-def* by *simp*

lemma *oor3I2*:

$\Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor3-def* by *simp*

lemma *oor3I3*:

$\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor3 \Delta \Delta_2 \Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor3-def* by *simp*

definition *oor4* **where**

$oor4 \Delta \Delta_2 \Delta_3 \Delta_4 \equiv \lambda w w1 w2 s1 s2 statA sv1 sv2 statO.$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$

lemma *oor4I1*:

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I2*:

$\Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I3*:

$\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

lemma *oor4I4*:

$\Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor4 \Delta \Delta_2 \Delta_3 \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor4-def* by *simp*

definition *oor5* where

$oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 \equiv \lambda w w1 w2 s1 s2 statA sv1 sv2 statO.$

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \vee$
 $\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \vee \Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO$
 \vee

$\Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

lemma *oor5I1*:

$\Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I2*:

$\Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I3*:

$\Delta_3 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I4*:

$\Delta_4 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *oor5I5*:

$\Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO \implies oor5 \Delta \Delta_2 \Delta_3 \Delta_4 \Delta_5 w w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *oor5-def* by *simp*

lemma *isIntO-match1*: $isIntO s1 \implies match1 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *match1-def* by *auto*

lemma *isIntO-match2*: $isIntO s2 \implies match2 \Delta w1 w2 s1 s2 statA sv1 sv2 statO$

unfolding *match2-def* by *auto*

lemma *isIntO-match*:

```

assumes  $\langle isIntO\ s1 \rangle$  and  $\langle isIntO\ s2 \rangle$ 
  and  $\langle match12\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \rangle$ 
  shows  $\langle react\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \rangle$ 
unfolding react-def apply (intro conjI)
subgoal
  using assms(1) by (rule isIntO-match1)
subgoal
  using assms(2) by (rule isIntO-match2)
subgoal
  using assms(3) by assumption
.

```

```

lemma match1-1-oorI1:
 $match1-1\ \Delta\ w1\ w2\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \implies$ 
 $match1-1\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO$ 
apply(rule match1-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match1-1-oorI2:
 $match1-1\ \Delta_2\ w1\ w2\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO \implies$ 
 $match1-1\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s1'\ s2\ statA\ sv1\ sv2\ statO$ 
apply(rule match1-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match1-oorI1:
 $match1\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$ 
 $match1\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
apply(rule match1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match1-oorI2:
 $match1\ \Delta_2\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$ 
 $match1\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 
apply(rule match1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match2-1-oorI1:
 $match2-1\ \Delta\ w1\ w2\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO \implies$ 
 $match2-1\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO$ 
apply(rule match2-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match2-1-oorI2:
 $match2-1\ \Delta_2\ w1\ w2\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO \implies$ 
 $match2-1\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s2\ s2'\ statA\ sv1\ sv2\ statO$ 
apply(rule match2-1-mono) unfolding le-fun-def oor-def by auto

```

```

lemma match2-oorI1:
 $match2\ \Delta\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO \implies$ 
 $match2\ (oor\ \Delta\ \Delta_2)\ w1\ w2\ s1\ s2\ statA\ sv1\ sv2\ statO$ 

```


apply(rule match2-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match2-oorI2:

match2 Δ_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies

match2 (oor Δ Δ_2) w1 w2 s1 s2 statA sv1 sv2 statO

apply(rule match2-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-oorI1:

match12 Δ w1 w2 s1 s2 statA sv1 sv2 statO \implies

match12 (oor Δ Δ_2) w1 w2 s1 s2 statA sv1 sv2 statO

apply(rule match12-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-oorI2:

match12 Δ_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies

match12 (oor Δ Δ_2) w1 w2 s1 s2 statA sv1 sv2 statO

apply(rule match12-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-1-oorI1:

match12-1 Δ w1 w2 s1' s2' statA' sv1 sv2 statO \implies

match12-1 (oor Δ Δ_2) w1 w2 s1' s2' statA' sv1 sv2 statO

apply(rule match12-1-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-1-oorI2:

match12-1 Δ_2 w1 w2 s1' s2' statA' sv1 sv2 statO \implies

match12-1 (oor Δ Δ_2) w1 w2 s1' s2' statA' sv1 sv2 statO

apply(rule match12-1-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-2-oorI1:

match12-2 Δ w1 w2 s2 s2' statA' sv1 sv2 statO \implies

match12-2 (oor Δ Δ_2) w1 w2 s2 s2' statA' sv1 sv2 statO

apply(rule match12-2-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-2-oorI2:

match12-2 Δ_2 w1 w2 s2 s2' statA' sv1 sv2 statO \implies

match12-2 (oor Δ Δ_2) w1 w2 s2 s2' statA' sv1 sv2 statO

apply(rule match12-2-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-12-oorI1:

match12-12 Δ w1 w2 s1' s2' statA' sv1 sv2 statO \implies

match12-12 (oor Δ Δ_2) w1 w2 s1' s2' statA' sv1 sv2 statO

apply(rule match12-12-mono) **unfolding** le-fun-def oor-def **by** auto

lemma match12-12-oorI2:

match12-12 Δ_2 w1 w2 s1' s2' statA' sv1 sv2 statO \implies

match12-12 (oor Δ Δ_2) w1 w2 s1' s2' statA' sv1 sv2 statO

apply(rule match12-12-mono) **unfolding** le-fun-def oor-def **by** auto

lemma *match-oorI1*:
 $react \Delta w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $react (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *match-oorI2*:
 $react \Delta_2 w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $react (oor \Delta \Delta_2) w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *match-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI1*:
 $proact \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $proact (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *proact-oorI2*:
 $proact \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $proact (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *proact-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-1-oorI1*:
 $move-1 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-1 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-1-oorI2*:
 $move-1 \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-1 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-1-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-2-oorI1*:
 $move-2 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-2 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-2-oorI2*:
 $move-2 \Delta_2 w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-2 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-2-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-12-oorI1*:
 $move-12 \Delta w w1 w2 s1 s2 statA sv1 sv2 statO \implies$
 $move-12 (oor \Delta \Delta_2) w w1 w2 s1 s2 statA sv1 sv2 statO$
apply(rule *move-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*

lemma *move-12-oorI2*:
move-12 Δ_2 w $w1$ $w2$ $s1$ $s2$ $statA$ $sv1$ $sv2$ $statO$ \implies
move-12 (*oor* Δ Δ_2) w $w1$ $w2$ $s1$ $s2$ $statA$ $sv1$ $sv2$ $statO$
apply(*rule* *move-12-mono*) **unfolding** *le-fun-def oor-def* **by** *auto*
end

context *Relative-Security-Determ*
begin

lemma *match1-1-defD*: *match1-1* Δ $w1$ $w2$ $s1$ $s1'$ $s2$ $statA$ $sv1$ $sv2$ $statO$ \longleftrightarrow
 \neg *finalV* $sv1$ \wedge $\Delta \infty$ $w1$ $w2$ $s1'$ $s2$ $statA$ (*nextO* $sv1$) $sv2$ $statO$
unfolding *match1-1-def validTrans-iff-next* **by** *simp*

lemma *match1-12-defD*: *match1-12* Δ $w1$ $w2$ $s1$ $s1'$ $s2$ $statA$ $sv1$ $sv2$ $statO$ \longleftrightarrow
 \neg *finalV* $sv1$ \wedge \neg *finalV* $sv2$ \wedge
 $\Delta \infty$ $w1$ $w2$ $s1'$ $s2$ $statA$ (*nextO* $sv1$) (*nextO* $sv2$) (*sstatO'* $statO$ $sv1$ $sv2$)
unfolding *match1-12-def validTrans-iff-next* **by** *simp*

lemmas *match1-defsD* = *match1-def match1-1-defD match1-12-defD*

lemma *match2-1-defD*: *match2-1* Δ $w1$ $w2$ $s1$ $s2$ $s2'$ $statA$ $sv1$ $sv2$ $statO$ \longleftrightarrow
 \neg *finalV* $sv2$ \wedge $\Delta \infty$ $w1$ $w2$ $s1$ $s2'$ $statA$ $sv1$ (*nextO* $sv2$) $statO$
unfolding *match2-1-def validTrans-iff-next* **by** *simp*

lemma *match2-12-defD*: *match2-12* Δ $w1$ $w2$ $s1$ $s2$ $s2'$ $statA$ $sv1$ $sv2$ $statO$ \longleftrightarrow
 \neg *finalV* $sv1$ \wedge \neg *finalV* $sv2$ \wedge $\Delta \infty$ $w1$ $w2$ $s1$ $s2'$ $statA$ (*nextO* $sv1$) (*nextO* $sv2$)
(*sstatO'* $statO$ $sv1$ $sv2$)
unfolding *match2-12-def validTrans-iff-next* **by** *simp*

lemmas *match2-defsD* = *match2-def match2-1-defD match2-12-defD*

lemma *match12-1-defD*: *match12-1* Δ $w1$ $w2$ $s1'$ $s2'$ $statA'$ $sv1$ $sv2$ $statO$ \longleftrightarrow
 \neg *finalV* $sv1$ \wedge $\Delta \infty$ $w1$ $w2$ $s1'$ $s2'$ $statA'$ (*nextO* $sv1$) $sv2$ $statO$
unfolding *match12-1-def validTrans-iff-next* **by** *simp*

lemma *match12-2-defD*: *match12-2* Δ $w1$ $w2$ $s1'$ $s2'$ $statA'$ $sv1$ $sv2$ $statO$ \longleftrightarrow
 \neg *finalV* $sv2$ \wedge $\Delta \infty$ $w1$ $w2$ $s1'$ $s2'$ $statA'$ $sv1$ (*nextO* $sv2$) $statO$
unfolding *match12-2-def validTrans-iff-next* **by** *simp*

lemma *match12-12-defD*: *match12-12* Δ $w1$ $w2$ $s1'$ $s2'$ $statA'$ $sv1$ $sv2$ $statO$ \longleftrightarrow
(*let* $statO' = sstatO' statO sv1 sv2$ *in*

$\neg \text{finalV } sv1 \wedge \neg \text{finalV } sv2 \wedge$
 $(\text{statA}' = \text{Diff} \longrightarrow \text{statO}' = \text{Diff}) \wedge$
 $\Delta \infty w1 w2 s1' s2' \text{statA}' (\text{nextO } sv1) (\text{nextO } sv2) \text{statO}'$
unfolding *match12-12-def validTrans-iff-next by simp*

lemmas *match12-defsD = match12-def match12-1-defD match12-2-defD match12-12-defD*

lemmas *match-deep-defsD = match1-defsD match2-defsD match12-defsD*

lemma *move-1-defD*: $\text{move-1 } \Delta w w1 w2 s1 s2 \text{statA } sv1 sv2 \text{statO} \longleftrightarrow$
 $\neg \text{finalV } sv1 \wedge \Delta w w1 w2 s1 s2 \text{statA } (\text{nextO } sv1) sv2 \text{statO}$
unfolding *move-1-def validTrans-iff-next by simp*

lemma *move-2-defD*: $\text{move-2 } \Delta w w1 w2 s1 s2 \text{statA } sv1 sv2 \text{statO} \longleftrightarrow$
 $\neg \text{finalV } sv2 \wedge \Delta w w1 w2 s1 s2 \text{statA } sv1 (\text{nextO } sv2) \text{statO}$
unfolding *move-2-def validTrans-iff-next by simp*

lemma *move-12-defD*: $\text{move-12 } \Delta w w1 w2 s1 s2 \text{statA } sv1 sv2 \text{statO} \longleftrightarrow$
 $(\text{let } \text{statO}' = \text{sstatO}' \text{statO } sv1 sv2 \text{ in}$
 $\neg \text{finalV } sv1 \wedge \neg \text{finalV } sv2 \wedge$
 $\Delta w w1 w2 s1 s2 \text{statA } (\text{nextO } sv1) (\text{nextO } sv2) \text{statO}')$
unfolding *move-12-def validTrans-iff-next by simp*

lemmas *proact-defsD = proact-def move-1-defD move-2-defD move-12-defD*

end

end

References

- [1] A. P. Brijesh Dongol, Matt Griffin and J. Wright. Relative security: Formally modeling and (dis)proving resilience against semantic optimization vulnerabilities. In *37th IEEE Computer Security Foundations Symposium, CSF 2024*. To appear.