

Regular Algebras

Simon Foster and Georg Struth

March 17, 2025

Abstract

Regular algebras axiomatise the equational theory of regular expressions as induced by regular language identity. We use Isabelle/HOL for a detailed systematic study of regular algebras given by Boffa, Conway, Kozen and Salomaa. We investigate the relationships between these classes, formalise a soundness proof for the smallest class (Salomaa's) and obtain completeness of the largest one (Boffa's) relative to a deep result by Krob. In addition we provide a large collection of regular identities in the general setting of Boffa's axiom.

Contents

1	Introductory Remarks	2
2	Diodoids, Powers and Finite Sums	2
3	Regular Algebras	4
3.1	Conway's Classical Axioms	4
3.2	Boffa's Axioms	5
3.3	Boffa Monoid Identities	10
3.4	Conway's Conjectures	11
3.5	Kozen's Kleene Algebras	13
3.6	Salomaa's Axioms	15
4	Models of Regular Algebra	16
4.1	Language Model of Salomaa Algebra	17
4.2	Regular Language Model of Salomaa Algebra	19
5	Pratt's Counterexamples	20
6	Variants of Regular Algebra	24

1 Introductory Remarks

These Isabelle theories complement the article on *On the Fine-Structure of Regular Algebra* [5]. For an introduction to the topic, conceptual explanations and references we refer to this article. Our regular algebra hierarchy is orthogonal to the Kleene algebra hierarchy in the Archive of Formal Proofs [1]; we have not aimed at an integration for pragmatic reasons.

2 Diods, Powers and Finite Sums

```
theory Diod-Power-Sum
  imports Kleene-Algebra.Diod Kleene-Algebra.Finite-Suprema
```

```
begin
```

We add a few facts about powers and finite sums—in fact, finite suprema—to an existing theory field for dioids.

```
context diod-one-zero
```

```
begin
```

```
lemma add-iso-r:  $y \leq z \implies x + y \leq x + z$ 
  ⟨proof⟩
```

```
notation power (⟨-⟩ [101,50] 100)
```

```
lemma power-subdist:  $x^n \leq (x + y)^n$ 
  ⟨proof⟩
```

```
lemma power-inductl-var:  $x \cdot y \leq y \implies x^n \cdot y \leq y$ 
  ⟨proof⟩
```

```
lemma power-inductr-var:  $y \cdot x \leq y \implies y \cdot x^n \leq y$ 
  ⟨proof⟩
```

```
definition powsum :: 'a ⇒ nat ⇒ nat ⇒ 'a (⟨-⟩ [101,50,50] 100) where
  powsum x m n = sum ((↑ x) {m..n + m})
```

```
lemmas powsum-simps = powsum-def atLeastAtMostSuc-conv numerals
```

```
lemma powsum1 [simp]:  $x_n^0 = x^n$ 
  ⟨proof⟩
```

```
lemma powsum2:  $x_n^{Suc m} = x_n^m + x^{n+Suc m}$ 
  ⟨proof⟩
```

```
lemma powsum-00 [simp]:  $x_0^0 = 1$ 
  ⟨proof⟩
```

```

lemma powsum-01 [simp]:  $x_0^1 = 1 + x$ 
  ⟨proof⟩

lemma powsum-10 [simp]:  $x_1^0 = x$ 
  ⟨proof⟩

lemma powsum-split:  $x_m^{i+\text{Suc } n} = x_m^i + x_{m+\text{Suc } i}^n$ 
  ⟨proof⟩

lemma powsum-split-var1:  $x_0^{n+1} = 1 + x_1^n$ 
  ⟨proof⟩

lemma powsum-split-var2 [simp]:  $x^m + x_0^m = x_0^m$ 
  ⟨proof⟩

lemma powsum-split-var3:  $x_0^{m+\text{Suc } n} = x_0^m + x_{0+\text{Suc } m}^n$ 
  ⟨proof⟩

lemma powsum-split-var4 [simp]:  $x_0^{m+n} + x_m^n = x_0^{m+n}$ 
  ⟨proof⟩

lemma powsum-split-var6:  $x_0^{(\text{Suc } k)+\text{Suc } n} = x_0^{\text{Suc } k} + x_{0+\text{Suc } (\text{Suc } k)}^n$ 
  ⟨proof⟩

lemma powsum-ext:  $x \leq x_0^{\text{Suc } n}$ 
  ⟨proof⟩

lemma powsum-one:  $1 \leq x_0^{\text{Suc } n}$ 
  ⟨proof⟩

lemma powsum-shift1:  $x \cdot x_m^n = x_{m+1}^n$ 
  ⟨proof⟩

lemma powsum-shift:  $x^k \cdot x_m^n = x_{k+m}^n$ 
  ⟨proof⟩

lemma powsum-prod-suc:  $x_0^m \cdot x_0^{\text{Suc } n} = x_0^{\text{Suc } (m+n)}$ 
  ⟨proof⟩

lemma powsum-prod:  $x_0^m \cdot x_0^n = x_0^{m+n}$ 
  ⟨proof⟩

end

end

```

3 Regular Algebras

```
theory Regular-Algebras
imports Diod-Power-Sum Kleene-Algebra Finite-Suprema Kleene-Algebra Kleene-Algebra
begin
```

3.1 Conway's Classical Axioms

Conway's classical axiomatisation of Regular Algebra from [4].

```
class star-diodoid = dioid-one-zero + star-op + plus-ord
```

```
class conway-diodoid = star-diodoid +
assumes C11:  $(x + y)^* = (x^* \cdot y)^* \cdot x^*$ 
and C12:  $(x \cdot y)^* = 1 + x \cdot (y \cdot x)^* \cdot y$ 
```

```
class strong-conway-diodoid = conway-diodoid +
assumes C13:  $(x^*)^* = x^*$ 
```

```
class C-algebra = strong-conway-diodoid +
assumes C14:  $x^* = (x^{n+1})^* \cdot x_0^n$ 
```

We tried to dualise using sublocales, but this causes an infinite loop on dual.dual.dual....

```
lemma (in conway-diodoid) C11-var:  $(x + y)^* = x^* \cdot (y \cdot x^*)^*$ 
⟨proof⟩
```

```
lemma (in conway-diodoid) dual-conway-diodoid:
class.conway-diodoid (+) (⊖) 1 0 (≤) (<) star
⟨proof⟩
```

```
lemma (in strong-conway-diodoid) dual-strong-conway-diodoid: class.strong-conway-diodoid
((+) ) ((⊖) ) 1 0 (≤) (<) star
⟨proof⟩
```

Nitpick finds counterexamples to the following claims.

```
lemma (in conway-diodoid) 1^* = 1
nitpick [expect=genuine] — 3-element counterexample
⟨proof⟩
```

```
lemma (in conway-diodoid) (x^*)^* = x^*
nitpick [expect=genuine] — 3-element counterexample
⟨proof⟩
```

```
context C-algebra
begin
```

```
lemma C-unfoldl [simp]: 1 + x · x^* = x^*
⟨proof⟩
```

```

lemma C-slide:  $(x \cdot y)^* \cdot x = x \cdot (y \cdot x)^*$ 
⟨proof⟩

lemma powsum-ub:  $i \leq n \implies x^i \leq x_0^n$ 
⟨proof⟩

lemma C14-aux:  $m \leq n \implies x^m \cdot (x^n)^* = (x^n)^* \cdot x^m$ 
⟨proof⟩

end

context dioid-one-zero
begin

lemma opp-power-def:
  power.power 1 (⊖) x n = xn
⟨proof⟩

lemma opp-powsum-def:
  dioid-one-zero.powsum (+) (⊖) 1 0 x m n = xmn
⟨proof⟩

end

lemma C14-dual:
  fixes x:'a::C-algebra
  shows x* = x0n · (xn+1)*
⟨proof⟩

lemma C-algebra: class.C-algebra (+) (⊖) (1:'a::C-algebra) 0 (≤) (<) star
⟨proof⟩

```

3.2 Boffa's Axioms

Boffa's two axiomatisations of Regular Algebra from [2, 3].

```

class B1-algebra = conway-dioid +
  assumes R:  $x \cdot x = x \implies x^* = 1 + x$ 

class B2-algebra = star-dioid +
  assumes B21:  $1 + x \leq x^*$ 
  and B22 [simp]:  $x^* \cdot x^* = x^*$ 
  and B23:  $\llbracket 1 + x \leq y; y \cdot y = y \rrbracket \implies x^* \leq y$ 

lemma (in B1-algebra) B1-algebra:
  class.B1-algebra (+) (⊖) 1 0 (≤) (<) star
⟨proof⟩

lemma (in B2-algebra) B2-algebra:

```

*class.B2-algebra (+) (\odot) 1 0 (\leq) ($<$) star
 $\langle proof \rangle$*

instance *B1-algebra \subseteq B2-algebra*
 $\langle proof \rangle$

context *B2-algebra*
begin

lemma *star-ref*: $1 \leq x^*$
 $\langle proof \rangle$

lemma *star-plus-one [simp]*: $1 + x^* = x^*$
 $\langle proof \rangle$

lemma *star-trans*: $x^* \cdot x^* \leq x^*$
 $\langle proof \rangle$

lemma *star-trans-eq [simp]*: $x^* \cdot x^* = x^*$
 $\langle proof \rangle$

lemma *star-invol [simp]*: $(x^*)^* = x^*$
 $\langle proof \rangle$

lemma *star-1l*: $x \cdot x^* \leq x^*$
 $\langle proof \rangle$

lemma *star-one [simp]*: $1^* = 1$
 $\langle proof \rangle$

lemma *star-subdist*: $x^* \leq (x + y)^*$
 $\langle proof \rangle$

lemma *star-iso*: $x \leq y \implies x^* \leq y^*$
 $\langle proof \rangle$

lemma *star2*: $(1 + x)^* = x^*$
 $\langle proof \rangle$

lemma *star-unfoldl*: $1 + x \cdot x^* \leq x^*$
 $\langle proof \rangle$

lemma *star-unfoldr*: $1 + x^* \cdot x \leq x^*$
 $\langle proof \rangle$

lemma *star-ext*: $x \leq x^*$
 $\langle proof \rangle$

lemma *star-1r*: $x^* \cdot x \leq x^*$

$\langle proof \rangle$

lemma star-unfoldl-eq [simp]: $1 + x \cdot x^* = x^*$
 $\langle proof \rangle$

lemma star-unfoldr-eq [simp]: $1 + x^* \cdot x = x^*$
 $\langle proof \rangle$

lemma star-prod-unfold-le: $(x \cdot y)^* \leq 1 + x \cdot (y \cdot x)^* \cdot y$
 $\langle proof \rangle$

lemma star-prod-unfold [simp]: $1 + x \cdot (y \cdot x)^* \cdot y = (x \cdot y)^*$
 $\langle proof \rangle$

lemma star-slide1: $(x \cdot y)^* \cdot x \leq x \cdot (y \cdot x)^*$
 $\langle proof \rangle$

lemma star-slide-var1: $x^* \cdot x \leq x \cdot x^*$
 $\langle proof \rangle$

lemma star-slide: $(x \cdot y)^* \cdot x = x \cdot (y \cdot x)^*$
 $\langle proof \rangle$

lemma star rtc1: $1 + x + x^* \cdot x^* \leq x^*$
 $\langle proof \rangle$

lemma star rtc1-eq: $1 + x + x^* \cdot x^* = x^*$
 $\langle proof \rangle$

lemma star-subdist-var-1: $x \leq (x + y)^*$
 $\langle proof \rangle$

lemma star-subdist-var-2: $x \cdot y \leq (x + y)^*$
 $\langle proof \rangle$

lemma star-subdist-var-3: $x^* \cdot y^* \leq (x + y)^*$
 $\langle proof \rangle$

lemma R-lemma:
 assumes $x \cdot x = x$
 shows $x^* = 1 + x$
 $\langle proof \rangle$

lemma star-denest-var-0: $(x + y)^* = (x^* \cdot y)^* \cdot x^*$
 $\langle proof \rangle$

lemma star-denest: $(x + y)^* = (x^* \cdot y^*)^*$
 $\langle proof \rangle$

lemma *star-sum-var*: $(x + y)^* = (x^* + y^*)^*$
⟨proof⟩

lemma *star-denest-var*: $(x + y)^* = x^* \cdot (y \cdot x^*)^*$
⟨proof⟩

lemma *star-denest-var-2*: $(x^* \cdot y^*)^* = x^* \cdot (y \cdot x^*)^*$
⟨proof⟩

lemma *star-denest-var-3*: $(x^* \cdot y^*)^* = x^* \cdot (y^* \cdot x^*)^*$
⟨proof⟩

lemma *star-denest-var-4*: $(x^* \cdot y^*)^* = (y^* \cdot x^*)^*$
⟨proof⟩

lemma *star-denest-var-5*: $x^* \cdot (y \cdot x^*)^* = y^* \cdot (x \cdot y^*)^*$
⟨proof⟩

lemma *star-denest-var-6*: $(x + y)^* = x^* \cdot y^* \cdot (x + y)^*$
⟨proof⟩

lemma *star-denest-var-7*: $(x + y)^* = (x + y)^* \cdot x^* \cdot y^*$
⟨proof⟩

lemma *star-denest-var-8*: $(x^* \cdot y^*)^* = x^* \cdot y^* \cdot (x^* \cdot y^*)^*$
⟨proof⟩

lemma *star-denest-var-9*: $(x^* \cdot y^*)^* = (x^* \cdot y^*)^* \cdot x^* \cdot y^*$
⟨proof⟩

lemma *star-slide-var*: $x^* \cdot x = x \cdot x^*$
⟨proof⟩

lemma *star-sum-unfold*: $(x + y)^* = x^* + x^* \cdot y \cdot (x + y)^*$
⟨proof⟩

lemma *troeger*: $x^* \cdot (y \cdot ((x + y)^* \cdot z) + z) = (x + y)^* \cdot z$
⟨proof⟩

lemma *meyer-1*: $x^* = (1 + x) \cdot (x \cdot x)^*$
⟨proof⟩

lemma *star-zero [simp]*: $0^* = 1$
⟨proof⟩

lemma *star-subsum [simp]*: $x^* + x^* \cdot x = x^*$
⟨proof⟩

lemma *prod-star-closure*: $x \leq z^* \implies y \leq z^* \implies x \cdot y \leq z^*$

```

⟨proof⟩

end

sublocale B2-algebra ⊆ B1-algebra
⟨proof⟩

context B2-algebra
begin

lemma power-le-star:  $x^n \leq x^*$ 
⟨proof⟩

lemma star-power-slide:
assumes  $k \leq n$ 
shows  $x^k \cdot (x^n)^* = (x^n)^* \cdot x^k$ 
⟨proof⟩

lemma powsum-le-star:  $x_m^n \leq x^*$ 
⟨proof⟩

lemma star-sum-power-slide:
assumes  $m \leq n$ 
shows  $x_0^m \cdot (x^n)^* = (x^n)^* \cdot x_0^m$ 
⟨proof⟩

lemma aarden-aux:
assumes  $y \leq y \cdot x + z$ 
shows  $y \leq y \cdot x^{(Suc\ n)} + z \cdot x^*$ 
⟨proof⟩

lemma conway-powerstar1:  $(x^{n+1})^* \cdot x_0^n \cdot (x^{n+1})^* \cdot x_0^n = (x^{n+1})^* \cdot x_0^n$ 
⟨proof⟩

lemma conway-powerstar2:  $1 + x \leq (x^{n+1})^* \cdot x_0^n$ 
⟨proof⟩

theorem powerstar:  $x^* = (x^{n+1})^* \cdot x_0^n$ 
⟨proof⟩

end

sublocale B2-algebra ⊆ strong-conway-diodoid
⟨proof⟩

sublocale B2-algebra ⊆ C-algebra
⟨proof⟩

```

The following fact could neither be verified nor falsified in Isabelle. It does

not hold for other reasons.

lemma (in *C-algebra*) $x \cdot x = x \longrightarrow x^* = 1 + x$
 $\langle proof \rangle$

3.3 Boffa Monoid Identities

typedef ('*a*, '*b*) *boffa-mon* = {*f* :: '*a*::{finite,monoid-mult} \Rightarrow '*b*::*B1-algebra*.
True}
 $\langle proof \rangle$

notation

Rep-boffa-mon ($\langle - \rangle$)

lemma *finite* (range (*Rep-boffa-mon M*))
 $\langle proof \rangle$

abbreviation *boffa-pair* :: ('*a*, '*b*) *boffa-mon* \Rightarrow '*a*::{finite,monoid-mult} \Rightarrow '*a* \Rightarrow
'*b*::*B1-algebra* **where**
 $boffa-pair\ i\ j \equiv \sum \{ x_k \mid k. i \cdot k = j \}$

notation

boffa-pair ($\langle -, - \rangle$)

abbreviation *conway-assms* **where**

conway-assms *x* \equiv ($\forall i\ j. (x_i \cdot x_j \leq x_{i \cdot j}) \wedge (x_{i,i})^* = x_{i,i})$

lemma *pair-one*: $x_{1,1} = x_1$
 $\langle proof \rangle$

definition *conway-assm1* **where** *conway-assm1* *x* = ($\forall i\ j. x_i \cdot x_j \leq x_{i \cdot j}$)
definition *conway-assm2* **where** *conway-assm2* *x* = ($\forall i. x_{i,i}^* = x_{i,i}$)

lemma *pair-star*:
assumes *conway-assm2* *x*
shows $x_1^* = x_1$
 $\langle proof \rangle$

lemma *conway-monoid-one*:
assumes *conway-assm2* *x*
shows $x_1 = 1 + x_1$
 $\langle proof \rangle$

lemma *conway-monoid-split*:
assumes *conway-assm2* *x*
shows $\sum \{x_i \mid i . i \in UNIV\} = 1 + \sum \{x_i \mid i . i \in UNIV\}$
 $\langle proof \rangle$

lemma *boffa-mon-aux1*: $\{x_{i,j} \mid i . j . i \in UNIV \wedge j \in UNIV\} = \{x_i \mid i . i \in UNIV\}$
 $\langle proof \rangle$

```

lemma sum-intro' [intro]:
   $\llbracket \text{finite } (A :: 'a::join-semilattice-zero\ set); \text{finite } B; \forall a \in A. \exists b \in B. a \leq b \rrbracket \implies$ 
 $\sum A \leq \sum B$ 
   $\langle proof \rangle$ 

lemma boffa-aux2:
   $\text{conway-assm1 } x \implies$ 
 $\sum \{x_i \cdot x_j \mid i \in UNIV \wedge j \in UNIV\} \leq \sum \{x_{i,j} \mid i \in UNIV \wedge j \in UNIV\}$ 
   $\langle proof \rangle$ 

lemma boffa-aux3:
  assumes conway-assm1 x
  shows  $(\sum \{x_i \mid i \in UNIV\}) + (\sum \{x_i \cdot x_j \mid i \in UNIV \wedge j \in UNIV\}) = (\sum \{x_i \mid i \in UNIV\})$ 
   $\langle proof \rangle$ 

lemma conway-monoid-identity:
  assumes conway-assm1 x conway-assm2 x
  shows  $(\sum \{x_i \mid i \in UNIV\})^* = (\sum \{x_i \mid i \in UNIV\})$ 
   $\langle proof \rangle$ 

```

3.4 Conway's Conjectures

```

class C0-algebra = strong-conway-diodid +
  assumes C0:  $x \cdot y = y \cdot z \implies x^* \cdot y = y \cdot z^*$ 

class C1l-algebra = strong-conway-diodid +
  assumes C1l:  $x \cdot y \leq y \cdot z \implies x^* \cdot y \leq y \cdot z^*$ 

class C1r-algebra = strong-conway-diodid +
  assumes C1r:  $y \cdot x \leq z \cdot y \implies y \cdot x^* \leq z^* \cdot y$ 

class C2l-algebra = conway-diodid +
  assumes C2l:  $x = y \cdot x \implies x = y^* \cdot x$ 

class C2r-algebra = conway-diodid +
  assumes C2r:  $x = x \cdot y \implies x = x \cdot y^*$ 

class C3l-algebra = conway-diodid +
  assumes C3l:  $x \cdot y \leq y \implies x^* \cdot y \leq y$ 

class C3r-algebra = conway-diodid +
  assumes C3r:  $y \cdot x \leq y \implies y \cdot x^* \leq y$ 

sublocale C1r-algebra  $\subseteq$  dual: C1l-algebra
  (+) ( $\odot$ ) 1 0 ( $\leq$ ) ( $<$ ) star
   $\langle proof \rangle$ 

```

sublocale $C2r\text{-algebra} \subseteq \text{dual: } C2l\text{-algebra}$
 $(+) (\odot) 1 0 (\leq) (<) \text{star}$
 $\langle proof \rangle$

sublocale $C3r\text{-algebra} \subseteq \text{dual: } C3l\text{-algebra}$
 $(+) (\odot) 1 0 (\leq) (<) \text{star}$
 $\langle proof \rangle$

lemma (in $C3l\text{-algebra}$) $k2\text{-var: } z + x \cdot y \leq y \implies x^* \cdot z \leq y$
 $\langle proof \rangle$

instance $C2l\text{-algebra} \subseteq B1\text{-algebra}$
 $\langle proof \rangle$

instance $C2r\text{-algebra} \subseteq B1\text{-algebra}$
 $\langle proof \rangle$

The following claims are refuted by Nitpick

lemma (in *conway-diodoid*)
assumes $x \cdot y = y \cdot z \implies x^* \cdot y = y \cdot z^*$
shows $1^* = 1$

$\langle proof \rangle$

lemma (in *conway-diodoid*)
assumes $x \cdot y \leq y \cdot z \implies x^* \cdot y \leq y \cdot z^*$
shows $1^* = 1$

$\langle proof \rangle$

The following fact could not be refuted by Nitpick or Quickcheck; but an infinite counterexample exists.

lemma (in $B1\text{-algebra}$) $x = x \cdot y \longrightarrow x = x \cdot y^*$
 $\langle proof \rangle$

instance $C3l\text{-algebra} \subseteq C2l\text{-algebra}$
 $\langle proof \rangle$

sublocale $C2l\text{-algebra} \subseteq C3l\text{-algebra}$
 $\langle proof \rangle$

sublocale $C1l\text{-algebra} \subseteq C3l\text{-algebra}$
 $\langle proof \rangle$

sublocale $C3l\text{-algebra} \subseteq C1l\text{-algebra}$
 $\langle proof \rangle$

sublocale $C1l\text{-algebra} \subseteq C2l\text{-algebra}$
 $\langle proof \rangle$

```

sublocale C3r-algebra ⊆ C2r-algebra
⟨proof⟩

sublocale C2r-algebra ⊆ C3r-algebra
⟨proof⟩

sublocale C1r-algebra ⊆ C3r-algebra
⟨proof⟩

sublocale C3r-algebra ⊆ C1r-algebra
⟨proof⟩

class C1-algebra = C1l-algebra + C1r-algebra

class C2-algebra = C2l-algebra + C2r-algebra

class C3-algebra = C3l-algebra + C3r-algebra

sublocale C0-algebra ⊆ C2-algebra
⟨proof⟩

sublocale C2-algebra ⊆ C0-algebra
⟨proof⟩

sublocale C2-algebra ⊆ C1-algebra ⟨proof⟩

sublocale C1-algebra ⊆ C2-algebra ⟨proof⟩

sublocale C2-algebra ⊆ C3-algebra ⟨proof⟩

sublocale C3-algebra ⊆ C2-algebra ⟨proof⟩

```

3.5 Kozen's Kleene Algebras

Kozen's Kleene Algebras [7, 6].

```

class Kl-base = star-diodoid +
assumes Kl:  $1 + x \cdot x^* \leq x^*$ 

class Kr-base = star-diodoid +
assumes Kr:  $1 + x^* \cdot x \leq x^*$ 

class K1l-algebra = Kl-base +
assumes star-inductl:  $x \cdot y \leq y \implies x^* \cdot y \leq y$ 

class K1r-algebra = Kr-base +
assumes star-inductr:  $y \cdot x \leq y \implies y \cdot x^* \leq y$ 

class K2l-algebra = Kl-base +

```

```

assumes star-inductl-var:  $z + x \cdot y \leq y \implies x^* \cdot z \leq y$ 

class K2r-algebra = Kr-base +
assumes star-inductr-var:  $z + y \cdot x \leq y \implies z \cdot x^* \leq y$ 

class K1-algebra = K1l-algebra + K1r-algebra

class K2-algebra = K2l-algebra + K2r-algebra

sublocale K1r-algebra  $\subseteq$  dual: K1l-algebra
  (+) ( $\odot$ ) 1 0 ( $\leq$ ) ( $<$ ) star
  ⟨proof⟩

sublocale K1l-algebra  $\subseteq$  B2-algebra
  ⟨proof⟩

sublocale K1r-algebra  $\subseteq$  B2-algebra
  ⟨proof⟩

sublocale K1l-algebra  $\subseteq$  C2l-algebra
  ⟨proof⟩

sublocale C2l-algebra  $\subseteq$  K1l-algebra
  ⟨proof⟩

sublocale K1r-algebra  $\subseteq$  C2r-algebra
  ⟨proof⟩

sublocale C2r-algebra  $\subseteq$  K1r-algebra
  ⟨proof⟩

sublocale K1-algebra  $\subseteq$  C0-algebra
  ⟨proof⟩

sublocale C0-algebra  $\subseteq$  K1l-algebra ⟨proof⟩

sublocale K2r-algebra  $\subseteq$  dual: K2l-algebra
  (+) ( $\odot$ ) 1 0 ( $\leq$ ) ( $<$ ) star
  ⟨proof⟩

sublocale K1l-algebra  $\subseteq$  K2l-algebra
  ⟨proof⟩

sublocale K2l-algebra  $\subseteq$  K1l-algebra
  ⟨proof⟩

sublocale K1r-algebra  $\subseteq$  K2r-algebra
  ⟨proof⟩

```

```

sublocale K2r-algebra ⊆ K1r-algebra
  ⟨proof⟩

sublocale kleene-algebra ⊆ K1-algebra
  ⟨proof⟩

sublocale K1-algebra ⊆ K2-algebra ⟨proof⟩

sublocale K2-algebra ⊆ koz: kleene-algebra
  ⟨proof⟩

```

3.6 Salomaa's Axioms

Salomaa's axiomatisations of Regular Algebra [9].

```

class salomaa-base = star-diodid +
  fixes ewp :: 'a ⇒ bool
  assumes S11:  $(1 + x)^* = x^*$ 
  and EWP : ewp x ←→ ( $\exists y. x = 1 + y \wedge \neg ewp y$ )

class Sr-algebra = salomaa-base +
  assumes S12r:  $1 + x^* \cdot x = x^*$ 
  and Ar : [ $\neg ewp y; x = x \cdot y + z$ ] ⇒⇒  $x = z \cdot y^*$ 

```

The following claim is ruled out by Nitpick. The unfold law cannot be weakened as in Kleene algebra.

```

lemma (in salomaa-base)
  assumes S12r':  $1 + x^* \cdot x \leq x^*$ 
  and Ar' : [ $\neg ewp y; x = x \cdot y + z$ ] ⇒⇒  $x = z \cdot y^*$ 
  shows  $x^* \leq 1 + x^* \cdot x$ 

```

⟨proof⟩

```

class Sl-algebra = salomaa-base +
  assumes S12l:  $1 + x \cdot x^* = x^*$ 
  and Al : [ $\neg ewp y; x = y \cdot x + z$ ] ⇒⇒  $x = y^* \cdot z$ 

```

```

class S-algebra = Sl-algebra + Sr-algebra

```

```

sublocale Sl-algebra ⊆ dual: Sr-algebra
  (+) (⊖) 1 0 (≤) (<) star eup
  ⟨proof⟩

```

```

context Sr-algebra
begin

```

```

lemma kozen-induct-r:
  assumes  $y \cdot x + z \leq y$ 
  shows  $z \cdot x^* \leq y$ 

```

```

⟨proof⟩

end

context Sl-algebra
begin

lemma kozen-induct-l:
  assumes  $x \cdot y + z \leq y$ 
  shows  $x^* \cdot z \leq y$ 
  ⟨proof⟩

end

sublocale Sr-algebra ⊆ K2r-algebra
  ⟨proof⟩

sublocale Sr-algebra ⊆ K1r-algebra ⟨proof⟩

sublocale Sl-algebra ⊆ K2l-algebra
  ⟨proof⟩

sublocale Sl-algebra ⊆ K1l-algebra ⟨proof⟩

sublocale S-algebra ⊆ K1-algebra ⟨proof⟩

sublocale S-algebra ⊆ K2-algebra ⟨proof⟩

The following claim could be refuted.

lemma (in K2-algebra) ( $\neg 1 \leq x$ ) —>  $x = x \cdot y + z \longrightarrow x = z \cdot y^*$ 
  ⟨proof⟩

class salomaa-conj-r = salomaa-base +
  assumes salomaa-small-unfold:  $1 + x^* \cdot x = x^*$ 
  assumes salomaa-small-r:  $\llbracket \neg \text{ewp } y ; x = x \cdot y + 1 \rrbracket \implies x = y^*$ 

sublocale Sr-algebra ⊆ salomaa-conj-r
  ⟨proof⟩

lemma (in salomaa-conj-r) ( $\neg \text{ewp } y$ )  $\wedge (x = x \cdot y + z) \longrightarrow x = z \cdot y^*$ 
  ⟨proof⟩

end

```

4 Models of Regular Algebra

```

theory Regular-Algebra-Models
  imports Regular-Algebras Kleene-Algebra.Kleene-Algebra-Models

```

begin

4.1 Language Model of Salomaa Algebra

abbreviation $w\text{-length} :: 'a \text{ list} \Rightarrow \text{nat} (\langle|\cdot|\rangle)$
where $|x| \equiv \text{length } x$

definition $l\text{-ewp} :: 'a \text{ lan} \Rightarrow \text{bool}$ **where**
 $l\text{-ewp } X \longleftrightarrow \{\}\subseteq X$

interpretation $\text{lan-kozen}: K2\text{-algebra} (+) (\cdot) 1 :: 'a \text{ lan } 0 \ (\subseteq) \ (\subset) \ \text{star} \ \langle\text{proof}\rangle$

interpretation $\text{lan-boffa}: B1\text{-algebra} (+) (\cdot) 1 :: 'a \text{ lan } 0 \ (\subseteq) \ (\subset) \ \text{star} \ \langle\text{proof}\rangle$

lemma $\text{length-lang-pow-lb}:$
assumes $\forall x \in X. |x| \geq k \ x \in X^{\wedge n}$
shows $|x| \geq k * n$
 $\langle\text{proof}\rangle$

lemma $l\text{-prod-elim}: w \in X \cdot Y \longleftrightarrow (\exists u v. w = u @ v \wedge u \in X \wedge v \in Y)$
 $\langle\text{proof}\rangle$

lemma $\text{power-minor-var}:$
assumes $\forall w \in X. k \leq |w|$
shows $\forall w \in X^{\text{Suc } n}. n * k \leq |w|$
 $\langle\text{proof}\rangle$

lemma $\text{power-lb}: (\forall w \in X. k \leq |w|) \longrightarrow (\forall w. w \in X^{\text{Suc } n} \longrightarrow n * k \leq |w|)$
 $\langle\text{proof}\rangle$

lemma $\text{prod-lb}:$
 $\llbracket (\forall w \in X. m \leq \text{length } w); (\forall w \in Y. n \leq \text{length } w) \rrbracket \Longrightarrow (\forall w \in (X \cdot Y). (m + n) \leq \text{length } w)$
 $\langle\text{proof}\rangle$

lemma $\text{suicide-aux-l}:$
 $\llbracket (\forall w \in Y. 0 \leq \text{length } w); (\forall w \in X^{\text{Suc } n}. n \leq \text{length } w) \rrbracket \Longrightarrow (\forall w \in X^{\text{Suc } n} \cdot Y. n \leq \text{length } w)$
 $\langle\text{proof}\rangle$

lemma $\text{suicide-aux-r}:$
 $\llbracket (\forall w \in Y. 0 \leq \text{length } w); (\forall w \in X^{\text{Suc } n}. n \leq \text{length } w) \rrbracket \Longrightarrow (\forall w \in Y \cdot X^{\text{Suc } n}. n \leq \text{length } w)$
 $\langle\text{proof}\rangle$

lemma $\text{word-suicide-l}:$
assumes $\neg l\text{-ewp } X \ Y \neq \{\}$
shows $(\forall w \in Y. \exists n. w \notin X^{\text{Suc } n} \cdot Y)$
 $\langle\text{proof}\rangle$

```

lemma word-suicide-r:
  assumes  $\neg l\text{-ewp } X \ Y \neq \{\}$ 
  shows  $(\forall w \in Y. \exists n. w \notin Y \cdot X^{\text{Suc } n})$ 
   $\langle proof \rangle$ 

lemma word-suicide-lang-l:  $\llbracket \neg l\text{-ewp } X; Y \neq \{\} \rrbracket \implies \exists n. \neg (Y \leq X^{\text{Suc } n} \cdot Y)$ 
   $\langle proof \rangle$ 

lemma word-suicide-lang-r:  $\llbracket \neg l\text{-ewp } X; Y \neq \{\} \rrbracket \implies \exists n. \neg (Y \leq Y \cdot X^{\text{Suc } n})$ 
   $\langle proof \rangle$ 

```

These duality results cannot be relocated easily

```

context K1-algebra
begin

```

```

lemma power-dual-transfer [simp]:
  power.power (1::'a) ( $\odot$ ) x n =  $x^n$ 
   $\langle proof \rangle$ 

```

```

lemma aarden-aux-l:
   $y \leq x \cdot y + z \implies y \leq x^{\text{Suc } n} \cdot y + x^* \cdot z$ 
   $\langle proof \rangle$ 

```

```
end
```

```

lemma arden-l:
  assumes  $\neg l\text{-ewp } y \ x = y \cdot x + z$ 
  shows  $x = y^* \cdot z$ 
   $\langle proof \rangle$ 

```

```

lemma arden-r:
  assumes  $\neg l\text{-ewp } y \ x = x \cdot y + z$ 
  shows  $x = z \cdot y^*$ 
   $\langle proof \rangle$ 

```

The following two facts provide counterexamples to Arden's rule if the empty word property is not considered.

```

lemma arden-l-counter:  $\exists (x::'a \ lan) (y::'a \ lan) (z::'a \ lan). x = y \cdot x + z \wedge x \neq y^* \cdot z$ 
   $\langle proof \rangle$ 

```

```

lemma arden-r-counter:  $\exists (x::'a \ lan) (y::'a \ lan) (z::'a \ lan). x = x \cdot y + z \wedge x \neq z \cdot y^*$ 
   $\langle proof \rangle$ 

```

```

interpretation lan-salomaa-l: Sl-algebra (+) ( $\cdot$ ) 1 :: 'a lan 0 ( $\subseteq$ ) ( $\subset$ ) star l-ewp
   $\langle proof \rangle$ 

```

interpretation *lan-salomaa-r*: *Sr-algebra* $(+)$ (\cdot) $1 :: 'a lan 0 (\subseteq) (\subset) star l-ewp \langle proof \rangle$

4.2 Regular Language Model of Salomaa Algebra

notation

Atom $(\langle \langle - \rangle \rangle)$ **and**
Plus (**infixl** $\langle +_r \rangle$ 65) **and**
Times (**infixl** $\langle \cdot_r \rangle$ 70) **and**
Star $(\langle -^*_r \rangle [101] 100)$ **and**
Zero $(\langle 0_r \rangle)$ **and**
One $(\langle 1_r \rangle)$

```
fun rexp-ewp :: 'a rexp  $\Rightarrow$  bool where
  rexp-ewp  $0_r = False$  |
  rexp-ewp  $1_r = True$  |
  rexp-ewp  $\langle a \rangle = False$  |
  rexp-ewp  $(s +_r t) = (\text{rexp-ewp } s \vee \text{rexp-ewp } t)$  |
  rexp-ewp  $(s \cdot_r t) = (\text{rexp-ewp } s \wedge \text{rexp-ewp } t)$  |
  rexp-ewp  $(s^*_r) = True$ 
```

abbreviation *ro*(*s*) \equiv (*if* (*rexp-ewp* *s*) *then* 1_r *else* 0_r)

lift-definition *r-ewp* :: '*a* *reg-lan* \Rightarrow bool **is** *l-ewp* $\langle proof \rangle$

lift-definition *r-lang* :: '*a* *rexp* \Rightarrow '*a* *reg-lan* **is** *lang*
 $\langle proof \rangle$

abbreviation *r-sim* :: '*a* *rexp* \Rightarrow '*a* *rexp* \Rightarrow bool (**infix** $\langle \sim \rangle$ 50) **where**
 $p \sim q \equiv r\text{-lang } p = r\text{-lang } q$

```
declare Rep-reg-lan [simp]
declare Rep-reg-lan-inverse [simp]
declare Abs-reg-lan-inverse [simp]
```

lemma *rexp-ewp-l-ewp*: *l-ewp* (*lang* *x*) $=$ *rexp-ewp* *x*
 $\langle proof \rangle$

theorem *regexp-ewp*:
defines *P-def*: *P(t)* $\equiv \exists t'. t \sim ro(t) +_r t' \wedge ro(t') = 0_r$
shows *P t*
 $\langle proof \rangle$

instantiation *reg-lan* :: (*type*) *Sr-algebra*
begin

lift-definition *ewp-reg-lan* :: '*a* *reg-lan* \Rightarrow bool **is** *l-ewp* $\langle proof \rangle$

instance $\langle proof \rangle$

```

end

instantiation reg-lan :: (type) Sl-algebra
begin

instance ⟨proof⟩
end

instance reg-lan :: (type) S-algebra ⟨proof⟩

theorem arden-regexp-l:
assumes ro(y) = 0r x ~ y ·r x +r z
shows x ~ y*r ·r z
⟨proof⟩

theorem arden-regexp-r:
assumes ro(y) = 0r x ~ x ·r y +r z
shows x ~ z ·r y*r
⟨proof⟩

end

```

5 Pratt's Counterexamples

```

theory Pratts-Counterexamples
  imports Regular-Algebras
begin

```

We create two regular algebra models due to Pratt [8] which are used to distinguish K1 algebras from K1l and K1r algebras.

```

datatype pratt1 =
  P1Bot (⟨⊥1⟩) |
  P1Nat nat (⟨[-]1⟩) |
  P1Infty (⟨∞1⟩) |
  P1Top (⟨⊤1⟩)

datatype pratt2 =
  P2Bot (⟨⊥2⟩) |
  P2Nat nat (⟨[-]2⟩) |
  P2Infty (⟨∞2⟩) |
  P2Top (⟨⊤2⟩)

fun pratt1-max where
  pratt1-max [x]1 [y]1 = [max x y]1 |
  pratt1-max x ⊥1 = x |
  pratt1-max ⊥1 y = y |
  pratt1-max ∞1 [y]1 = ∞1 |
  pratt1-max [y]1 ∞1 = ∞1 |
  pratt1-max ∞1 ∞1 = ∞1 |

```

```

pratt1-max  $\top_1$   $x = \top_1$  |
pratt1-max  $x \top_1 = \top_1$ 

fun pratt1-plus :: pratt1  $\Rightarrow$  pratt1  $\Rightarrow$  pratt1 (infixl  $\langle +_1 \rangle$  65) where
   $[x]_1 +_1 [y]_1 = [x + y]_1$  |
   $\perp_1 +_1 x = \perp_1$  |
   $x +_1 \perp_1 = \perp_1$  |
   $\infty_1 +_1 [0]_1 = \infty_1$  |
   $[0]_1 +_1 \infty_1 = \infty_1$  |
   $\infty_1 +_1 \infty_1 = \top_1$  |
   $\top_1 +_1 x = \top_1$  |
   $x +_1 \top_1 = \top_1$  |
   $[x]_1 +_1 \infty_1 = \infty_1$  |
   $\infty_1 +_1 x = \top_1$ 

lemma plusl-top-infty:  $\top_1 +_1 \infty_1 = \top_1$ 
   $\langle proof \rangle$ 

lemma plusl-infty-top:  $\infty_1 +_1 \top_1 = \top_1$ 
   $\langle proof \rangle$ 

lemma plusl-bot-infty:  $\perp_1 +_1 \infty_1 = \perp_1$ 
   $\langle proof \rangle$ 

lemma plusl-infty-bot:  $\infty_1 +_1 \perp_1 = \perp_1$ 
   $\langle proof \rangle$ 

lemma plusl-zero-infty:  $[0]_1 +_1 \infty_1 = \infty_1$ 
   $\langle proof \rangle$ 

lemma plusl-infty-zero:  $\infty_1 +_1 [0]_1 = \infty_1$ 
   $\langle proof \rangle$ 

lemma plusl-infty-num [simp]:  $x > 0 \implies \infty_1 +_1 [x]_1 = \top_1$ 
   $\langle proof \rangle$ 

lemma plusl-num-infty [simp]:  $x > 0 \implies [x]_1 +_1 \infty_1 = \infty_1$ 
   $\langle proof \rangle$ 

fun pratt2-max where
  pratt2-max  $[x]_2 [y]_2 = [\max x y]_2$  |
  pratt2-max  $x \perp_2 = x$  |
  pratt2-max  $\perp_2 y = y$  |
  pratt2-max  $\infty_2 [y]_2 = \infty_2$  |
  pratt2-max  $[y]_2 \infty_2 = \infty_2$  |
  pratt2-max  $\infty_2 \infty_2 = \infty_2$  |
  pratt2-max  $\top_2 x = \top_2$  |
  pratt2-max  $x \top_2 = \top_2$ 

```

```

fun pratt2-plus :: pratt2  $\Rightarrow$  pratt2  $\Rightarrow$  pratt2 (infixl  $\langle +_2 \rangle$  65) where
   $[x]_2 +_2 [y]_2 = [x + y]_2 \mid$ 
   $\perp_2 +_2 x = \perp_2 \mid$ 
   $x +_2 \perp_2 = \perp_2 \mid$ 
   $\infty_2 +_2 [\theta]_2 = \infty_2 \mid$ 
   $[\theta]_2 +_2 \infty_2 = \infty_2 \mid$ 
   $\infty_2 +_2 \infty_2 = \top_2 \mid$ 
   $\top_2 +_2 x = \top_2 \mid$ 
   $x +_2 \top_2 = \top_2 \mid$ 
   $x +_2 \infty_2 = \top_2 \mid$ 
   $\infty_2 +_2 [x]_2 = \infty_2$ 

instantiation pratt1 :: selective-semiring
begin

  definition zero-pratt1-def:
     $0 \equiv \perp_1$ 

  definition one-pratt1-def:
     $1 \equiv [\theta]_1$ 

  definition plus-pratt1-def:
     $x + y \equiv \text{pratt1-max } x \ y$ 

  definition times-pratt1-def:
     $x * y \equiv x +_1 y$ 

  definition less-eq-pratt1-def:
     $(x::\text{pratt1}) \leq y \equiv x + y = y$ 

  definition less-pratt1-def:
     $(x::\text{pratt1}) < y \equiv x \leq y \wedge x \neq y$ 

  instance
   $\langle \text{proof} \rangle$ 
end

instance pratt1 :: diod-one-zero  $\langle \text{proof} \rangle$ 

instantiation pratt2 :: selective-semiring
begin

  definition zero-pratt2-def:
     $0 \equiv \perp_2$ 

  definition one-pratt2-def:
     $1 \equiv [\theta]_2$ 

  definition times-pratt2-def:

```

```

 $x * y \equiv x +_2 y$ 

definition plus-pratt2 :: pratt2  $\Rightarrow$  pratt2  $\Rightarrow$  pratt2 where
  plus-pratt2 x y  $\equiv$  pratt2-max x y

definition less-eq-pratt2-def:
   $(x::\text{pratt2}) \leq y \equiv x + y = y$ 

definition less-pratt2-def:
   $(x::\text{pratt2}) < y \equiv x \leq y \wedge x \neq y$ 

instance
   $\langle \text{proof} \rangle$ 
end

lemma top-greatest:  $x \leq \top_1$ 
   $\langle \text{proof} \rangle$ 

instantiation pratt1 :: star-op
begin

  definition star-pratt1 where
     $x^* \equiv \text{if } (x = \perp_1 \vee x = [0]_1) \text{ then } [0]_1 \text{ else } \top_1$ 
    instance  $\langle \text{proof} \rangle$ 
  end

  instantiation pratt2 :: star-op
  begin
    definition star-pratt2 where
       $x^* \equiv \text{if } (x = \perp_2 \vee x = [0]_2) \text{ then } [0]_2 \text{ else } \top_2$ 
    instance  $\langle \text{proof} \rangle$ 
  end

  instance pratt1 :: K1r-algebra
   $\langle \text{proof} \rangle$ 

  instance pratt2 :: K1l-algebra
   $\langle \text{proof} \rangle$ 

  lemma one-star-top:  $[1]_1^* = \top_1$ 
   $\langle \text{proof} \rangle$ 

  lemma pratt1-kozen-1l-counterexample:
   $\exists x y :: \text{pratt1}. \neg (x \cdot y \leq y \wedge x^* \cdot y \leq y)$ 
   $\langle \text{proof} \rangle$ 

  lemma pratt2-kozen-1r-counterexample:
   $\exists x y :: \text{pratt2}. \neg (y \cdot x \leq y \wedge y \cdot x^* \leq y)$ 
   $\langle \text{proof} \rangle$ 

```

end

6 Variants of Regular Algebra

```
theory Regular-Algebra-Variants
  imports Regular-Algebras Pratts-Counterexamples
begin
```

Replacing Kozen's induction axioms by Boffa's leads to incompleteness.

```
lemma (in star-dioïd)
  assumes  $\bigwedge x. 1 + x \cdot x^* = x^*$ 
  and  $\bigwedge x. 1 + x^* \cdot x = x^*$ 
  and  $\bigwedge x. x \cdot x = x \implies x^* = 1 + x$ 
  shows  $\bigwedge x y. (x + y)^* = x^* \cdot (y \cdot x^*)^*$ 
```

(proof)

```
lemma (in star-dioïd)
  assumes  $\bigwedge x. 1 + x \cdot x^* = x^*$ 
  and  $\bigwedge x. 1 + x^* \cdot x = x^*$ 
  and  $\bigwedge x. x \cdot x = x \implies x^* = 1 + x$ 
  shows  $\bigwedge x y. x \leq y \longrightarrow x^* \leq y^*$ 
```

(proof)

```
lemma (in star-dioïd)
  assumes  $\bigwedge x. 1 + x \cdot x^* = x^*$ 
  and  $\bigwedge x. 1 + x^* \cdot x = x^*$ 
  and  $\bigwedge x y. 1 + x \leq y \wedge y \cdot y \leq y \longrightarrow x^* \leq y$ 
  shows  $\bigwedge x. 1 + x \leq x^*$ 
(proof)
```

```
lemma (in star-dioïd)
  assumes  $\bigwedge x. 1 + x \cdot x^* = x^*$ 
  and  $\bigwedge x. 1 + x^* \cdot x = x^*$ 
  and  $\bigwedge x y. 1 + x \leq y \wedge y \cdot y \leq y \longrightarrow x^* \leq y$ 
  shows  $\bigwedge x. x^* = (1 + x)^*$ 
(proof)
```

```
lemma (in star-dioïd)
  assumes  $\bigwedge x. 1 + x \cdot x^* = x^*$ 
  and  $\bigwedge x. 1 + x^* \cdot x = x^*$ 
  and  $\bigwedge x y. 1 + x + y \cdot y \leq y \implies x^* \leq y$ 
  shows  $\bigwedge x. x^* \cdot x^* \leq x^*$ 
(proof)
```

```
lemma (in star-dioïd)
  assumes  $\bigwedge x. 1 + x \cdot x^* = x^*$ 
```

and $\bigwedge x. 1 + x^* \cdot x = x^*$
and $\bigwedge x y z. x \cdot y = y \cdot z \implies x^* \cdot y = y \cdot z^*$
shows $1^* = 1$

$\langle proof \rangle$

lemma (in star-diod)
assumes $\bigwedge x. 1 + x \cdot x^* = x^*$
and $\bigwedge x. 1 + x^* \cdot x = x^*$
and $\bigwedge x y z. x \cdot y \leq y \cdot z \implies x^* \cdot y \leq y \cdot z^*$
and $\bigwedge x y z. y \cdot x \leq z \cdot y \implies y \cdot x^* \leq z^* \cdot y$
shows $1^* = 1$

$\langle proof \rangle$

lemma (in star-diod)
assumes $\bigwedge x. 1 + x \cdot x^* = x^*$
and $\bigwedge x. 1 + x^* \cdot x = x^*$
and $\bigwedge x y. x = y \cdot x \implies x = y^* \cdot x$
and $\bigwedge x y. x = x \cdot y \implies x = x \cdot y^*$
shows $\bigwedge x y. y \cdot x \leq y \implies y \cdot x^* \leq y$

$\langle proof \rangle$

class $C3l\text{-var} = \text{conway-diod} +$
assumes $C3l\text{-var}: z + x \cdot y \leq y \implies x^* \cdot z \leq y$

class $C3r\text{-var} = \text{conway-diod} +$
assumes $C3r\text{-var}: z + y \cdot x \leq y \implies z \cdot x^* \leq y$

class $C3\text{-var} = C3l\text{-var} + C3r\text{-var}$

sublocale $C3l\text{-var} \subseteq C3l\text{-algebra}$
 $\langle proof \rangle$

sublocale $C3l\text{-algebra} \subseteq C3l\text{-var}$
 $\langle proof \rangle$

sublocale $C3\text{-var} \subseteq C3\text{-algebra}$
 $\langle proof \rangle$

sublocale $C3\text{-algebra} \subseteq C3\text{-var}$
 $\langle proof \rangle$

class $Brtc\text{-algebra} = \text{star-diod} +$
assumes $rtc1: 1 + x^* \cdot x^* + x \leq x^*$
and $rtc2: 1 + x + y \cdot y \leq y \implies x^* \leq y$

sublocale $B2\text{-algebra} \subseteq Brtc\text{-algebra}$
 $\langle proof \rangle$

```

sublocale Brtc-algebra  $\subseteq$  B2-algebra
⟨proof⟩

class wB1-algebra = conway-diodoid +
assumes wR:  $x \cdot x \leq x \implies x^* = 1 + x$ 

sublocale wB1-algebra  $\subseteq$  B1-algebra
⟨proof⟩

lemma (in B1-algebra) one-plus-star:  $x^* = (1 + x)^*$ 
⟨proof⟩

sublocale B1-algebra  $\subseteq$  wB1-algebra
⟨proof⟩

end

```

References

- [1] A. Armstrong, G. Struth, and T. Weber. Kleene algebra. *Archive of Formal Proofs*, 2013. http://isa-afp.org/entries/Kleene_Algebra.shtml, Formal proof development.
- [2] M. Boffa. Une remarque sur les systèmes complets d'identités rationnelles. *Informatique théorique et applications*, 24(4):419–423, 1990.
- [3] M. Boffa. Une condition impliquant toutes les identités rationnelles. *Informatique théorique et applications*, 29(6):515–518, 1995.
- [4] J. H. Conway. *Regular Algebra and Finite Machines*. Chapman and Hall, 1971.
- [5] S. Foster and G. Struth. On the fine-structure of regular algebra. *J. Automated Reasoning*, 54(2):165–197, 2015.
- [6] D. Kozen. On Kleene algebras and closed semirings. In B. Rovan, editor, *MFCS'90*, volume 452 of *LNCS*, pages 26–47. Springer, 1990.
- [7] D. Kozen. A completeness theorem for Kleene algebras and the algebra of regular events. *Information and Computation*, 110(2):366–390, 1994.
- [8] V. Pratt. Action logic and pure induction. Technical Report STAN-CS-90-1343, Department of Computer Science, Stanford University, 1990.
- [9] A. Salomaa. Two complete axiom systems for the algebra of regular events. *J. ACM*, 13(1):158–169, 1966.