

Ramsey's Theorem

Tom Ridge

October 11, 2017

Abstract

The infinite form of Ramsey's Theorem is proved following Boolos and Jeffrey, Chapter 26.

Contents

1 Ramsey's Theorem	1
1.1 Library lemmas	1
1.2 Dependent Choice Variant	2
1.3 Partitions	3
1.4 Ramsey's theorem	3

1 Ramsey's Theorem

```
theory Ramsey
imports Main HOL-Library.Infinite-Set
begin
```

```
declare [[simp-depth-limit = 5]]
```

1.1 Library lemmas

```
lemma infinite-inj-infinite-image: infinite Z ==> inj-on f Z ==> infinite (f ` Z)
  apply(rule ccontr)
  apply(simp)
  apply(force dest: finite-imageD)
done
```

```
lemma infinite-dom-finite-rng: [[ infinite A; finite (f ` A) ]] ==> ? b : f ` A.
infinite {a : A. f a = b}
  apply(rule ccontr) apply(simp)
  apply(subgoal-tac UNION A (% b. {a : A. f a = f b}) = A) prefer 2 apply(blast)
  apply(subgoal-tac (UN c : f ` A. {a : A. f a = c}) = (UN b:A. {a : A. f a = f
b})) prefer 2 apply(blast)
  apply(subgoal-tac finite (UN c:f ` A. {a : A. f a = c})) apply(force)
```

```

apply(rule finite-UN-I)
apply(auto)
done

```

```

lemma infinite-mem: infinite X ==> ? x. x : X
apply(rule ccontr)
apply(force)
done

```

```

lemma not-empty-least: (Y::nat set) ~ = {} ==> ? m. m : Y & (! m'. m' : Y
--> m <= m')
apply(rule-tac x=LEAST x. x : Y in exI)
apply(rule)
apply(rule LeastI-ex) apply(force)
apply(rule)
apply(rule)
apply(rule Least-le)
apply(assumption)
done

```

1.2 Dependent Choice Variant

```

primrec choice :: ('a => bool) => ('a => 'a => bool) => nat => 'a where
  choice P R 0 = (SOME x. P x)
| choice P R (Suc n) = (let x = choice P R n in SOME y. P y & R x y)

```

lemma dc:

```

(! x y z. R x y & R y z --> R x z)
& (? x0. P x0)
& (! x. P x --> (? y. P y & R x y))
--> (? f::nat=>'b. (! n. P (f n)) & (! n m. R (f n) (f (n+m+1))))

```

```

apply(intro allI impI, elim exE conjE)
apply(rule-tac x=choice P R in exI)
apply(subgoal-tac (ALL n. P (choice P R n))) prefer 2
apply(rule, induct-tac n)
apply(simp add: Let-def) apply(rule someI-ex) apply(blast)
apply(simp add: Let-def) apply(subgoal-tac P (SOME y. P y & R (choice P R
na) y) & R (choice P R na) (SOME y. P y & R (choice P R na) y)) apply(blast)
apply(rule someI-ex) apply(blast)
apply(rule) apply(assumption)

```

```

apply(subgoal-tac ! n. R (choice P R n) (choice P R (Suc n))) prefer 2
apply(rule)
apply(simp add: Let-def)
apply(subgoal-tac P (SOME y. P y & R (choice P R n) y) & R (choice P R n)
(SOME y. P y & R (choice P R n) y)) apply(blast)

```

```

apply(rule someI-ex) apply(force)

apply(rule) apply(rule)
apply(induct-tac m)
apply(force)
apply(drule-tac x=n+na+1 in spec) back
apply(force simp del: choice.simps)
done

```

1.3 Partitions

definition

```

part :: nat => nat => 'a set => ('a set => nat) => bool where
part r s Y f = (! X. X <= Y & finite X & card X = r --> f X < s)

```

```

lemma part: [| infinite YY; part (Suc n) s YY f; yy : YY |] ==> part n s (YY
- {yy}) (%u. f (insert yy u))
apply(simp add: part-def)
apply(intro allI impI, elim exE conjE)
apply(drule-tac x=insert yy X in spec)
apply(force simp: card-Diff-singleton-if)
done

```

```

lemma part-subset: part (Suc n) s YY f ==> Y <= YY ==> part (Suc n) s Y
f
apply(simp add: part-def)
apply(blast)
done

```

1.4 Ramsey's theorem

lemma ramsey:

```

! (s::nat) (r::nat) (YY::'a set) (f::'a set => nat).
infinite YY
& (! X. X <= YY & finite X & card X = r --> f X < s)
--> (? Y' t'.
Y' <= YY
& infinite Y'
& t' < s
& (! X. X <= Y' & finite X & card X = r --> f X = t'))
apply(simp add: part-def[symmetric] del: ex-simps)
apply(rule, rule, rule-tac nat.induct)

```

```

apply(intro allI impI)
apply(rule-tac x=YY in exI)
apply(rule-tac x=f {} in exI)
apply(force simp: part-def)

```

```

apply(intro allI impI) apply(elim exE conjE)

```

```

apply(subgoal-tac
  ? g.
  (! m::nat. let (y,Y,t) = (g m) in
    y : YY & y ~: Y
    & Y <= YY & infinite Y
    & t < s
    & (! X. X <= Y & finite X & card X = nat --> (f o insert y) X = t)
    & (! m m'.
      let (y,Y,t) = (g m) in
      let (y',Y',t') = (g (m+m'+1)) in
      y' : Y
      & Y' <= Y
    )
  )
)
prefer 2
apply(cut-tac

  P = % gn.
  let (y,Y,t) = (gn) in
  y : YY & y ~: Y
  & Y <= YY & infinite Y
  & t < s
  & (! X. X <= Y & finite X & card X = nat --> (f o insert y) X = t)

  and R = % gn gn'.
  let (y,Y,t) = (gn) in
  let (y',Y',t') = (gn') in
  y' : Y
  & Y' <= Y

  in dc)
apply(erule impE)

  apply(intro conjI)
  apply(intro allI impI, elim conjE) apply(simp add: Let-def split-beta) ap-
ply(blast)
  apply(subgoal-tac ? yy. yy : YY) prefer 2 apply(rule infinite-mem) ap-
ply(assumption)
  apply(elim exE conjE)
  apply(drule-tac x=YY - {yy} in spec) apply(drule-tac x=f o insert yy in
spec)
  apply(erule impE) apply(simp) apply(rule part) apply(assumption+)
  apply(elim exE conjE)
  apply(rule-tac x=(yy, Y', t') in exI) apply(simp) apply(blast)
  apply(intro allI impI)
  apply(case-tac x) apply(rename-tac yx b Yx tx)

```

apply(*subgoal-tac* ? $yx'. yx' : Yx$) **prefer** 2 **apply**(*rule infinite-mem*) **ap-**
ply(*force*)
apply(*elim exE conjE*)
apply(*drule-tac* $x=Yx - \{yx'\}$ **in** *spec*)
apply(*drule-tac* $x=f o$ **insert** yx' **in** *spec*)
apply(*erule impE*) **apply**(*simp*) **apply**(*elim exE conjE*) **apply**(*rule part*)
apply(*assumption+*)
apply(*rule part-subset*) **apply**(*assumption*) **apply**(*assumption*) **apply**(*assumption*)
apply(*elim exE conjE*)
apply(*rule-tac* $x=(yx', Y', t')$ **in** *exI*) **apply**(*simp*) **apply**(*blast*)
apply(*assumption*)

apply(*elim exE conjE*)

apply(*subgoal-tac* ? $s'. s' < s \ \& \ \text{infinite } \{n. (\% n. \text{let } (y, Y, t) = g \ n \ \text{in } t) \ n = s'\}$) **prefer** 2

apply(*subgoal-tac* ? $s' : ((\% n. \text{let } (Y, y, t) = g \ n \ \text{in } t) \ ' \ \text{UNIV}). \ \text{infinite } \{n : \text{UNIV}. (\text{let } (Y, y, t) = g \ n \ \text{in } t) = s'\}$) **prefer** 2
apply(*rule infinite-dom-finite-rng*) **apply**(*simp*)
apply(*simp* (*no-asm*) *add: finite-nat-iff-bounded*)
apply(*rule-tac* $x=s$ **in** *exI*)
apply(*rule*)
apply(*simp* *add: image-iff*) **apply**(*elim exE conjE*)
apply(*drule-tac* $x=xa$ **in** *spec*) **apply**(*force simp* *add: Let-def split-beta*)
apply(*elim bexE conjE*) **apply**(*rule-tac* $x=s'$ **in** *exI*) **apply**(*simp*)
apply(*simp* *add: image-iff*) **apply**(*elim exE conjE*) **apply**(*drule-tac* $x=x$ **in** *spec*) **apply**(*force simp: Let-def split-beta*)

apply(*elim exE conjE*)
apply(*rule-tac* $x=(\% n. \text{let } (y, Y, t) = g \ n \ \text{in } y) \ ' \ \{n. (\% n. \text{let } (y, Y, t) = g \ n \ \text{in } t) \ n = s'\}$ **in** *exI*)
apply(*rule-tac* $x=s'$ **in** *exI*)

apply(*subgoal-tac* *inj* ($\% n. \text{let } (y, Y, t) = g \ n \ \text{in } y$)) **prefer** 2
apply(*simp* *add: inj-on-def*)

apply(*subgoal-tac* *ALL* $x \ y. x < y \ \& \ (\text{let } (y, Y, t) = g \ x \ \text{in } y) = (\text{let } (y, Y, t) = g \ y \ \text{in } y) \ \text{-->} \ x = y$)
apply(*intro allI impI*)
apply(*subgoal-tac* $x < y \ | \ x = y \ | \ y < x$) **prefer** 2 **apply**(*arith*)
apply(*elim disjE*)
apply(*drule-tac* $x=x$ **in** *spec*) **back** **back** **apply**(*drule-tac* $x=y$ **in** *spec*) **back**
back **apply**(*force simp: Let-def*)
apply(*assumption*)
apply(*drule-tac* $x=y$ **in** *spec*) **back** **back** **apply**(*drule-tac* $x=x$ **in** *spec*) **back**
back **apply**(*force simp: Let-def*)
apply(*intro allI impI*)
apply(*drule-tac* $x=x$ **in** *spec*) **apply**(*drule-tac* $x=x$ **in** *spec*) **apply**(*drule-tac*

```

x=y-(Suc x) in spec) apply(force simp: Let-def)

  apply(intro allI conjI)

  apply(drule-tac x=xa in spec)
  apply(force simp add: Let-def split-beta)

  apply(rule infinite-inj-infinite-image) apply(assumption)
  apply(simp add: inj-on-def)

  apply(simp)

  apply(intro allI impI, elim exE conjE)
  apply(simp add: subset-image-iff) apply(elim exE conjE)
  apply(subgoal-tac ? a. a : AA & (! a'. a' : AA --> a <= a')) prefer 2
  apply(rule not-empty-least) apply(force)
  apply(elim exE conjE)
  apply(case-tac g a rule: prod.exhaust)
  apply(rename-tac b) apply(case-tac b rule: prod.exhaust) apply(rename-tac ya
  b Ya ta)
  apply(subgoal-tac ya : X) prefer 2 apply(force intro!: rev-image-eqI simp:
  Let-def)
  apply(drule-tac s=X in sym)
  apply(subgoal-tac f X = (f o insert ya) (X - {ya})) apply(simp)
  apply(drule-tac x=a in spec)
  apply(simp add: Let-def) apply(elim exE conjE)
  apply(drule-tac x=X-{ya} in spec) back apply(erule impE)
  apply(simp)
  apply(rule)
  apply(rule)
  apply(drule-tac t=X in sym) apply(simp)
  apply(simp add: image-iff) apply(elim bexE exE conjE) apply(rename-tac a')
  apply(subgoal-tac a' ~ = a) prefer 2 apply(force)
  apply(drule-tac x=a in spec)
  apply(drule-tac x=a'-Suc a in spec) back
  apply(simp add: Let-def split-beta) apply(case-tac g a' rule: prod.exhaust) ap-
  ply(case-tac ba rule: prod.exhaust) apply(rename-tac ya' ba Ya' ta') apply(simp
  add: Let-def split-beta)
  apply(drule-tac x=a' in spec) apply(erule impE) apply(force)
  apply(force)
  apply(force simp add: card-Diff-singleton-if)
  apply(subgoal-tac ta = s') apply(simp) apply(force)
  apply(simp) apply(rule-tac f=f in arg-cong) apply(force)
  done

end

```