

# A Sound and Complete Calculus for Probability Inequalities

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March 17, 2025

### Abstract

We give a sound and complete multiple-conclusion calculus  $\$ \vdash$  for finitely additive probability inequalities. In particular, we show

$$\sim\Gamma\$ \vdash \sim\Phi \equiv \forall \mathcal{P} \in \text{probabilities}. \sum \phi \leftarrow \Phi. \mathcal{P}\phi \leq \sum \gamma \leftarrow \Gamma. \mathcal{P}\gamma$$

... where  $\sim\Gamma$  is the negation of all of the formulae in  $\Gamma$  (and similarly for  $\sim\Phi$ ). We prove this by using an abstract form of *MaxSAT*. We also show

$$\text{MaxSAT}(\sim\Gamma @ \Phi) + c \leq \text{length } \Gamma \equiv \forall \mathcal{P} \in \text{probabilities}. \left( \sum \phi \leftarrow \Phi. \mathcal{P}\phi \right) + c \leq \sum \gamma \leftarrow \Gamma. \mathcal{P}\gamma$$

Finally, we establish a *collapse theorem*, which asserts that  $(\sum \phi \leftarrow \Phi. \mathcal{P}\phi) + c \leq \sum \gamma \leftarrow \Gamma. \mathcal{P}\gamma$  holds for all probabilities  $\mathcal{P}$  if and only if  $(\sum \phi \leftarrow \Phi. \delta\phi) + c \leq \sum \gamma \leftarrow \Gamma. \delta\gamma$  holds for all binary-valued probabilities  $\delta$ .

# Contents

<b>1</b>	<b>Introduction</b>	<b>2</b>
<b>2</b>	<b>Measure Deduction and Counting Deduction</b>	<b>4</b>
2.1	Definition of Measure Deduction . . . . .	4
2.2	Definition of the Stronger Theory Relation . . . . .	5
2.3	The Stronger Theory Relation is a Preorder . . . . .	6
2.4	The Stronger Theory Relation is a Subrelation of Measure Deduction . . . . .	10
2.5	Measure Deduction is a Preorder . . . . .	20
2.6	Measure Deduction Cancellation Rules . . . . .	86
2.7	Measure Deduction Substitution Rules . . . . .	93
2.8	Measure Deduction Sum Rules . . . . .	94
2.9	Measure Deduction Exchange Rule . . . . .	95
2.10	Definition of Counting Deduction . . . . .	96
2.11	Converting Back and Forth from Counting Deduction to Mea- sure Deduction . . . . .	96
2.12	Measure Deduction Soundness . . . . .	99
<b>3</b>	<b>MaxSAT</b>	<b>106</b>
3.1	Definition of Relative Maximal Clause Collections . . . . .	106
3.2	Definition of MaxSAT . . . . .	111
3.3	Reducing Counting Deduction to MaxSAT . . . . .	113
<b>4</b>	<b>Inequality Completeness For Probability Logic</b>	<b>156</b>
4.1	Limited Counting Deduction Completeness . . . . .	156
4.2	Measure Deduction Completeness . . . . .	159
4.3	Counting Deduction Completeness . . . . .	161
4.4	Collapse Theorem For Probability Logic . . . . .	162
4.5	MaxSAT Completeness For Probability Logic . . . . .	167

# Chapter 1

## Introduction

```
theory Probability-Inequality-Completeness
imports
  Suppes-Theorem.Probability-Logic
begin
```

```
unbundle no funcset-syntax
```

We introduce a novel logical calculus and prove completeness for probability inequalities. This is a vast generalization of *Suppes' Theorem* which lays the foundation for this theory.

We provide two new logical judgements: *measure deduction* ( $\$ \vdash$ ) and *counting deduction* ( $\# \vdash$ ). Both judgements capture a notion of *measure* or *quantity*. In both cases premises must be partially or completely *consumed* in sense to prove multiple conclusions. That is to say, a portion of the premises must be used to prove each conclusion which cannot be reused. Counting deduction counts the number of times a particular conclusion can be proved (as the name implies), while measure deduction includes multiple, different conclusions which must be proven via the premises.

We also introduce an abstract notion of MaxSAT, which is the maximal number of clauses in a list of clauses that can be simultaneously satisfied.

We show the following are equivalent:

- $\sim \Gamma \ \$ \vdash \sim \Phi$
- $(\sim \Gamma \ @ \ \Phi) \ \# \vdash \ (\text{length } \Phi) \perp$
- $\text{MaxSAT } (\sim \Gamma \ @ \ \Phi) \leq \text{length } \Gamma$
- $\forall \delta \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \delta \varphi) \leq (\sum \gamma \leftarrow \Gamma. \delta \gamma)$
- $\forall \mathcal{P} \in \text{probabilities. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$

In the special case of MaxSAT, we show the following are equivalent:

- $MaxSAT(\sim \Gamma @ \Phi) + c \leq length \Gamma$
- $\forall \delta \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \delta \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \delta \gamma)$
- $\forall \mathcal{P} \in \text{probabilities. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$

## Chapter 2

# Measure Deduction and Counting Deduction

### 2.1 Definition of Measure Deduction

To start, we introduce a common combinator for modifying functions that take two arguments.

**definition**  $uncurry :: ('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a \times 'b \Rightarrow 'c$   
**where**  $uncurry-def [simp]: uncurry f = (\lambda (x, y). f x y)$

Our new logical calculus is a recursively defined relation ( $\$ \vdash$ ) using *list deduction* ( $:\vdash$ ).

We call our new logical relation *measure deduction*:

**primrec** (in *classical-logic*)  
 $measure-deduction :: 'a\ list \Rightarrow 'a\ list \Rightarrow bool$  (**infix**  $\langle \$ \vdash \rangle$  60)  
**where**  
   $\Gamma \$ \vdash [] = True$   
   $|\ \Gamma \$ \vdash (\varphi \# \Phi) =$   
     $(\exists \Psi. mset (map\ snd\ \Psi) \subseteq\# mset\ \Gamma$   
       $\wedge map (uncurry (\sqcup)) \Psi :\vdash \varphi$   
       $\wedge map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus (map\ snd\ \Psi) \$ \vdash \Phi)$

Let us briefly analyze what the above definition is saying.

From the above we must find a special list-of-pairs  $\Psi$ , which we refer to as a *witness*, in order to establish  $\Gamma \$ \vdash \varphi \# \Phi$ .

We may motivate measure deduction as follows. In the simplest case we know  $\mathcal{P} \varphi \leq \mathcal{P} \psi + \Sigma$  if and only if  $\mathcal{P} (\chi \sqcup \varphi) + \mathcal{P} (\sim \chi \sqcup \varphi) \leq \mathcal{P} \psi + \Sigma$ , or equivalently  $\mathcal{P} (\chi \sqcup \varphi) + \mathcal{P} (\chi \rightarrow \varphi) \leq \mathcal{P} \psi + \Sigma$ . So it suffices to prove  $\mathcal{P} (\chi \sqcup \varphi) \leq \mathcal{P} \psi$  and  $\mathcal{P} (\chi \rightarrow \varphi) \leq \Sigma$ . Here  $[(\chi, \varphi)]$  is like the *witness* in our recursive definition, which reflects the  $\exists \Psi. \dots$  formula is our definition. The fact that measure deduction reflects proving theorems

in the theory of inequalities of probability logic is the elementary intuition behind the soundness theorem we will ultimately prove in §2.12.

A key difference from the simple motivation above is that, as in the case of Suppes' Theorem where we prove  $\sim \Gamma \vdash \sim \varphi$  if and only if  $\mathcal{P} \varphi \leq (\sum \gamma \leftarrow \Gamma . \mathcal{P} \gamma)$  for all  $\mathcal{P}$ , soundness in this context means  $\sim \Gamma \text{\$}\vdash \sim \Phi$  implies  $\forall \mathcal{P}. (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma) \geq (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi)$ .

Another way of thinking about measure deduction is to think of  $\Gamma$  and  $\Sigma$  as bags of balls of soft clay and  $\Gamma \text{\$}\vdash \Sigma$  meaning that we have shown  $\Gamma$  is *heavier than*  $\Sigma$  (ignoring, for the moment, that  $\text{\$}\vdash$  is not totally ordered). We have a scale ( $\vdash$ ) that lets us weigh several things on the left and one thing on the right at a time. We go through each clay ball  $\sigma$  in  $\Sigma$  one at a time without replacement, putting  $\sigma$  on the right of the scale. Then, we take a bunch of clay balls from  $\Gamma$ , cut them up as necessary (that is the  $\psi \sqcup \gamma$  trick using the witness  $\Psi$ ), and show they are heavier using our scale. We take the parts  $\psi \rightarrow \gamma$  that we didn't use and put them back in our bag  $\Gamma$ . We will be able to reuse them later. If we can do this trick for every element  $\sigma$  in  $\Sigma$  successively using combinations of split leftovers in  $\Gamma$ , then we can show  $\Gamma$  is heavier than  $\Sigma$  (i.e.,  $\Gamma \text{\$}\vdash \Sigma$ ).

## 2.2 Definition of the Stronger Theory Relation

We next turn to looking at a subrelation of  $\text{\$}\vdash$ , which we call the *stronger theory* relation ( $\preceq$ ). Here we construe a *theory* as a list of propositions. We say theory  $\Gamma$  is *stronger than*  $\Sigma$  where, for each element  $\sigma$  in  $\Sigma$ , we can take an element  $\gamma$  of  $\Gamma$  *without replacement* such that  $\vdash \gamma \rightarrow \sigma$ .

To motivate this notion, let's reuse the metaphor that  $\Gamma$  and  $\Sigma$  are bags of balls of clay, and we need to show  $\Gamma$  is heavier without simply weighing the two bags. A sufficient (but incomplete) approach is to take each ball of clay  $\sigma$  in  $\Sigma$  and find another ball of clay  $\gamma$  in  $\Gamma$  (without replacement) that is heavier. This simple approach avoids the complexity of iteratively cutting up balls of clay.

**definition** (in *implication-logic*)

*stronger-theory-relation* :: 'a list  $\Rightarrow$  'a list  $\Rightarrow$  bool (infix  $\langle \preceq \rangle$  100)

**where**

$$\begin{aligned} \Sigma \preceq \Gamma = & \\ & (\exists \Phi. \text{map snd } \Phi = \Sigma \\ & \wedge \text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } \Gamma \\ & \wedge (\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma)) \end{aligned}$$

**abbreviation** (in *implication-logic*)

*stronger-theory-relation-op* :: 'a list  $\Rightarrow$  'a list  $\Rightarrow$  bool (infix  $\langle \succeq \rangle$  100)

**where**

$$\Gamma \succeq \Sigma \equiv \Sigma \preceq \Gamma$$

## 2.3 The Stronger Theory Relation is a Preorder

Next, we show that  $(\preceq)$  is a preorder by establishing reflexivity and transitivity.

We first prove the following lemma with respect to multisets and stronger theories.

**lemma** (in *implication-logic*) *msub-stronger-theory-intro*:

**assumes**  $mset\ \Sigma \subseteq\# mset\ \Gamma$

**shows**  $\Sigma \preceq \Gamma$

**proof** –

**let**  $?\Delta\Sigma = map\ (\lambda\ x.\ (x,x))\ \Sigma$

**have**  $map\ snd\ ?\Delta\Sigma = \Sigma$

**by** (*induct*  $\Sigma$ , *simp*, *simp*)

**moreover have**  $map\ fst\ ?\Delta\Sigma = \Sigma$

**by** (*induct*  $\Sigma$ , *simp*, *simp*)

**hence**  $mset\ (map\ fst\ ?\Delta\Sigma) \subseteq\# mset\ \Gamma$

**using** *assms* **by** *simp*

**moreover have**  $\forall\ (\gamma,\sigma) \in set\ ?\Delta\Sigma.\ \vdash\ \gamma \rightarrow \sigma$

**by** (*induct*  $\Sigma$ , *simp*, *simp*,

*metis list-implication.simps(1) list-implication-axiom-k*)

**ultimately show** *?thesis* **using** *stronger-theory-relation-def* **by** (*simp*, *blast*)

**qed**

The *reflexive* property immediately follows:

**lemma** (in *implication-logic*) *stronger-theory-reflexive* [*simp*]:  $\Gamma \preceq \Gamma$

**using** *msub-stronger-theory-intro* **by** *auto*

**lemma** (in *implication-logic*) *weakest-theory* [*simp*]:  $[] \preceq \Gamma$

**using** *msub-stronger-theory-intro* **by** *auto*

**lemma** (in *implication-logic*) *stronger-theory-empty-list-intro* [*simp*]:

**assumes**  $\Gamma \preceq []$

**shows**  $\Gamma = []$

**using** *assms* *stronger-theory-relation-def* **by** *simp*

Next, we turn to proving transitivity. We first prove two permutation theorems.

**lemma** (in *implication-logic*) *stronger-theory-right-permutation*:

**assumes**  $\Gamma \rightleftharpoons \Delta$

**and**  $\Sigma \preceq \Gamma$

**shows**  $\Sigma \preceq \Delta$

**proof** –

**from** *assms(1)* **have**  $mset\ \Gamma = mset\ \Delta$

**by** *simp*

**thus** *?thesis*

**using** *assms(2)* *stronger-theory-relation-def*

**by** *fastforce*



qed

**lemma** (in *implication-logic*) *stronger-theory-left-permutation*:

assumes  $\Sigma \rightleftharpoons \Delta$

and  $\Sigma \preceq \Gamma$

shows  $\Delta \preceq \Gamma$

**proof** –

have  $\forall \Sigma \Gamma. \Sigma \rightleftharpoons \Delta \longrightarrow \Sigma \preceq \Gamma \longrightarrow \Delta \preceq \Gamma$

**proof** (*induct*  $\Delta$ )

case *Nil*

then show *?case* by *simp*

**next**

case (*Cons*  $\delta \Delta$ )

{

fix  $\Sigma \Gamma$

assume  $\Sigma \rightleftharpoons (\delta \# \Delta) \Sigma \preceq \Gamma$

from *this* obtain  $\Phi$  where  $\Phi$ :

*map snd*  $\Phi = \Sigma$

*mset (map fst*  $\Phi) \subseteq\#$  *mset*  $\Gamma$

$\forall (\gamma, \delta) \in \text{set } \Phi. \vdash \gamma \rightarrow \delta$

using *stronger-theory-relation-def* by *fastforce*

with  $\langle \Sigma \rightleftharpoons (\delta \# \Delta) \rangle$  have  $\delta \in\#$  *mset* (*map snd*  $\Phi$ )

by *fastforce*

from *this* obtain  $\gamma$  where  $\gamma: (\gamma, \delta) \in\#$  *mset*  $\Phi$

by (*induct*  $\Phi$ , *fastforce+*)

let  $?\Phi_0 = \text{remove1 } (\gamma, \delta) \Phi$

let  $?\Sigma_0 = \text{map snd } ?\Phi_0$

from  $\gamma \Phi(2)$  have *mset* (*map fst*  $?\Phi_0$ )  $\subseteq\#$  *mset* (*remove1*  $\gamma \Gamma$ )

by (*metis ex-mset*

*list-subtract-monotonic*

*list-subtract-mset-homomorphism*

*mset-remove1*

*remove1-pairs-list-projections-fst*)

moreover have *mset*  $?\Phi_0 \subseteq\#$  *mset*  $\Phi$  by *simp*

with  $\Phi(3)$  have  $\forall (\gamma, \delta) \in \text{set } ?\Phi_0. \vdash \gamma \rightarrow \delta$  by *fastforce*

ultimately have  $?\Sigma_0 \preceq \text{remove1 } \gamma \Gamma$

unfolding *stronger-theory-relation-def* by *blast*

moreover have  $\Delta \rightleftharpoons (\text{remove1 } \delta \Sigma)$  using  $\langle \Sigma \rightleftharpoons (\delta \# \Delta) \rangle$

by (*metis perm-remove-perm perm-sym remove-hd*)

moreover from  $\gamma \Phi(1)$  have *mset*  $?\Sigma_0 = \text{mset } (\text{remove1 } \delta \Sigma)$

using *remove1-pairs-list-projections-snd*

by *fastforce*

hence  $?\Sigma_0 \rightleftharpoons \text{remove1 } \delta \Sigma$

by *blast*

ultimately have  $\Delta \preceq \text{remove1 } \gamma \Gamma$  using *Cons*

by *presburger*

from *this* obtain  $\Psi_0$  where  $\Psi_0$ :

*map snd*  $\Psi_0 = \Delta$

*mset (map fst*  $\Psi_0) \subseteq\#$  *mset* (*remove1*  $\gamma \Gamma$ )

```

     $\forall (\gamma, \delta) \in \text{set } \Psi_0. \vdash \gamma \rightarrow \delta$ 
    using stronger-theory-relation-def by fastforce
  let ? $\Psi = (\gamma, \delta) \# \Psi_0$ 
  have map_snd ? $\Psi = (\delta \# \Delta)$ 
    by (simp add:  $\Psi_0(1)$ )
  moreover have mset (map fst ? $\Psi$ )  $\subseteq\#$  mset ( $\gamma \# (\text{remove1 } \gamma \Gamma)$ )
    using  $\Psi_0(2)$  by auto
  moreover from  $\gamma \Phi(3) \Psi_0(3)$  have  $\forall (\gamma, \sigma) \in \text{set } ?\Psi. \vdash \gamma \rightarrow \sigma$  by auto
  ultimately have  $(\delta \# \Delta) \preceq (\gamma \# (\text{remove1 } \gamma \Gamma))$ 
    unfolding stronger-theory-relation-def by metis
  moreover from  $\gamma \Phi(2)$  have  $\gamma \in\#$  mset  $\Gamma$ 
    using mset-subset-eqD by fastforce
  hence  $(\gamma \# (\text{remove1 } \gamma \Gamma)) \doteq \Gamma$ 
    by auto
  ultimately have  $(\delta \# \Delta) \preceq \Gamma$ 
    using stronger-theory-right-permutation by blast
}
then show ?case by blast
qed
with assms show ?thesis by blast
qed

```

lemma (in implication-logic) stronger-theory-transitive:

assumes  $\Sigma \preceq \Delta$  and  $\Delta \preceq \Gamma$   
 shows  $\Sigma \preceq \Gamma$

proof –

have  $\forall \Delta \Gamma. \Sigma \preceq \Delta \longrightarrow \Delta \preceq \Gamma \longrightarrow \Sigma \preceq \Gamma$

proof (induct  $\Sigma$ )

case Nil

then show ?case using stronger-theory-relation-def by simp

next

case (Cons  $\sigma \Sigma$ )

{

fix  $\Delta \Gamma$

assume  $(\sigma \# \Sigma) \preceq \Delta$   $\Delta \preceq \Gamma$

from this obtain  $\Phi$  where  $\Phi$ :

map\_snd  $\Phi = \sigma \# \Sigma$

mset (map fst  $\Phi$ )  $\subseteq\#$  mset  $\Delta$

$\forall (\delta, \sigma) \in \text{set } \Phi. \vdash \delta \rightarrow \sigma$

using stronger-theory-relation-def by (simp, metis)

let ? $\delta = \text{fst } (\text{hd } \Phi)$

from  $\Phi(1)$  have  $\Phi \neq []$  by (induct  $\Phi$ , simp+)

hence ? $\delta \in\#$  mset (map fst  $\Phi$ ) by (induct  $\Phi$ , simp+)

with  $\Phi(2)$  have ? $\delta \in\#$  mset  $\Delta$  by (meson mset-subset-eqD)

hence mset (map fst (remove1 (hd  $\Phi$ )  $\Phi$ ))  $\subseteq\#$  mset (remove1 ? $\delta \Delta$ )

using  $\langle \Phi \neq [] \rangle \Phi(2)$

by (simp,

metis

diff-single-eq-union

*hd-in-set*  
*image-mset-add-mset*  
*insert-subset-eq-iff*  
*set-mset-mset*

**moreover have**  $remove1 (hd \Phi) \Phi = tl \Phi$   
**using**  $\langle \Phi \neq [] \rangle$   
**by** (*induct*  $\Phi$ , *simp+*)

**moreover from**  $\Phi(1)$  **have**  $map\ snd (tl \Phi) = \Sigma$   
**by** (*simp add*: *map-tl*)

**moreover from**  $\Phi(3)$  **have**  $\forall (\delta, \sigma) \in set (tl \Phi). \vdash \delta \rightarrow \sigma$   
**by** (*simp add*:  $\langle \Phi \neq [] \rangle list.set-sel(2)$ )

**ultimately have**  $\Sigma \preceq remove1\ ?\delta \Delta$   
**using** *stronger-theory-relation-def* **by** *auto*

**from**  $\langle ?\delta \in \# mset \Delta \rangle$  **have**  $?\delta \# (remove1\ ?\delta \Delta) \equiv \Delta$   
**by** *fastforce*

**with**  $\langle \Delta \preceq \Gamma \rangle$  **have**  $(?\delta \# (remove1\ ?\delta \Delta)) \preceq \Gamma$   
**using** *stronger-theory-left-permutation perm-sym* **by** *blast*

**from this obtain**  $\Psi$  **where**  $\Psi$ :  
 $map\ snd \Psi = (?\delta \# (remove1\ ?\delta \Delta))$   
 $mset (map\ fst \Psi) \subseteq \# mset \Gamma$   
 $\forall (\gamma, \delta) \in set \Psi. \vdash \gamma \rightarrow \delta$   
**using** *stronger-theory-relation-def* **by** (*simp*, *metis*)

**let**  $?\gamma = fst (hd \Psi)$   
**from**  $\Psi(1)$  **have**  $\Psi \neq []$  **by** (*induct*  $\Psi$ , *simp+*)  
**hence**  $?\gamma \in \# mset (map\ fst \Psi)$  **by** (*induct*  $\Psi$ , *simp+*)  
**with**  $\Psi(2)$  **have**  $?\gamma \in \# mset \Gamma$  **by** (*meson mset-subset-eqD*)  
**hence**  $mset (map\ fst (remove1 (hd \Psi) \Psi)) \subseteq \# mset (remove1\ ?\gamma \Gamma)$   
**using**  $\langle \Psi \neq [] \rangle \Psi(2)$   
**by** (*simp*,  
*metis*  
*diff-single-eq-union*  
*hd-in-set*  
*image-mset-add-mset*  
*insert-subset-eq-iff*  
*set-mset-mset*)

**moreover from**  $\langle \Psi \neq [] \rangle$  **have**  $remove1 (hd \Psi) \Psi = tl \Psi$   
**by** (*induct*  $\Psi$ , *simp+*)

**moreover from**  $\Psi(1)$  **have**  $map\ snd (tl \Psi) = (remove1\ ?\delta \Delta)$   
**by** (*simp add*: *map-tl*)

**moreover from**  $\Psi(3)$  **have**  $\forall (\gamma, \delta) \in set (tl \Psi). \vdash \gamma \rightarrow \delta$   
**by** (*simp add*:  $\langle \Psi \neq [] \rangle list.set-sel(2)$ )

**ultimately have**  $remove1\ ?\delta \Delta \preceq remove1\ ?\gamma \Gamma$   
**using** *stronger-theory-relation-def* **by** *auto*

**with**  $\langle \Sigma \preceq remove1\ ?\delta \Delta \rangle$  *Cons.hyps* **have**  $\Sigma \preceq remove1\ ?\gamma \Gamma$   
**by** *blast*

**from this obtain**  $\Omega_0$  **where**  $\Omega_0$ :  
 $map\ snd \Omega_0 = \Sigma$   
 $mset (map\ fst \Omega_0) \subseteq \# mset (remove1\ ?\gamma \Gamma)$   
 $\forall (\gamma, \sigma) \in set \Omega_0. \vdash \gamma \rightarrow \sigma$

```

    using stronger-theory-relation-def by (simp, metis)
  let ?Ω = (?γ, σ) # Ω0
  from Ω0(1) have map_snd ?Ω = σ # Σ by simp
  moreover from Ω0(2) have mset (map fst ?Ω) ⊆# mset (?γ # (remove1
?γ Γ))
    by simp
  moreover from Φ(1) Ψ(1) have σ = snd (hd Φ) ?δ = snd (hd Ψ) by
fastforce+
  with Φ(3) Ψ(3) ⟨Φ ≠ []⟩ ⟨Ψ ≠ []⟩ hd-in-set have ⊢ ?δ → σ ⊢ ?γ → ?δ
    by fastforce+
  hence ⊢ ?γ → σ using modus-ponens hypothetical-syllogism by blast
  with Ω0(3) have ∀ (γ,σ) ∈ set ?Ω. ⊢ γ → σ
    by auto
  ultimately have (σ # Σ) ⪯ (?γ # (remove1 ?γ Γ))
    unfolding stronger-theory-relation-def
    by metis
  moreover from ⟨?γ ∈# mset Γ⟩ have (?γ # (remove1 ?γ Γ)) ⇒ Γ
    by force
  ultimately have (σ # Σ) ⪯ Γ
    using stronger-theory-right-permutation
    by blast
}
then show ?case by blast
qed
thus ?thesis using assms by blast
qed

```

## 2.4 The Stronger Theory Relation is a Subrelation of Measure Deduction

Next, we show that  $\Gamma \succeq \Sigma$  implies  $\Gamma \text{ \$}\vdash \Sigma$ . Before doing so we establish several helpful properties regarding the stronger theory relation ( $\succeq$ ).

**lemma** (in *implication-logic*) *stronger-theory-witness*:

```

  assumes σ ∈ set Σ
  shows Σ ⪯ Γ = (∃ γ ∈ set Γ. ⊢ γ → σ ∧ (remove1 σ Σ) ⪯ (remove1 γ Γ))
proof (rule iffI)
  assume Σ ⪯ Γ
  from this obtain Φ where Φ:
    map_snd Φ = Σ
    mset (map fst Φ) ⊆# mset Γ
    ∀ (γ,σ) ∈ set Φ. ⊢ γ → σ
  unfolding stronger-theory-relation-def by blast
  from assms Φ(1) obtain γ where γ: (γ, σ) ∈# mset Φ
  by (induct Φ, fastforce+)
  hence γ ∈# mset (map fst Φ) by force
  hence γ ∈# mset Γ using Φ(2)
  by (meson mset-subset-eqD)

```

**moreover**  
**let**  $?\Phi_0 = \text{remove1 } (\gamma, \sigma) \Phi$   
**let**  $? \Sigma_0 = \text{map snd } ?\Phi_0$   
**from**  $\gamma \Phi(2)$  **have**  $\text{mset } (\text{map fst } ?\Phi_0) \subseteq\# \text{mset } (\text{remove1 } \gamma \Gamma)$   
**by** (*metis*  
*ex-mset*  
*list-subtract-monotonic*  
*list-subtract-mset-homomorphism*  
*remove1-pairs-list-projections-fst*  
*mset-remove1*)  
**moreover have**  $\text{mset } ?\Phi_0 \subseteq\# \text{mset } \Phi$  **by** *simp*  
**with**  $\Phi(3)$  **have**  $\forall (\gamma, \sigma) \in \text{set } ?\Phi_0. \vdash \gamma \rightarrow \sigma$  **by** *fastforce*  
**ultimately have**  $? \Sigma_0 \preceq \text{remove1 } \gamma \Gamma$   
**unfolding** *stronger-theory-relation-def* **by** *blast*  
**moreover from**  $\gamma \Phi(1)$  **have**  $\text{mset } ?\Sigma_0 = \text{mset } (\text{remove1 } \sigma \Sigma)$   
**using** *remove1-pairs-list-projections-snd*  
**by** *fastforce*  
**hence**  $? \Sigma_0 \Rightarrow \text{remove1 } \sigma \Sigma$   
**by** *linarith*  
**ultimately have**  $\text{remove1 } \sigma \Sigma \preceq \text{remove1 } \gamma \Gamma$   
**using** *stronger-theory-left-permutation*  
**by** *blast*  
**moreover from**  $\gamma \Phi(3)$  **have**  $\vdash \gamma \rightarrow \sigma$  **by** (*simp, fast*)  
**moreover from**  $\gamma \Phi(2)$  **have**  $\gamma \in\# \text{mset } \Gamma$   
**using** *mset-subset-eqD* **by** *fastforce*  
**ultimately show**  $\exists \gamma \in \text{set } \Gamma. \vdash \gamma \rightarrow \sigma \wedge (\text{remove1 } \sigma \Sigma) \preceq (\text{remove1 } \gamma \Gamma)$  **by**  
*auto*  
**next**  
**assume**  $\exists \gamma \in \text{set } \Gamma. \vdash \gamma \rightarrow \sigma \wedge (\text{remove1 } \sigma \Sigma) \preceq (\text{remove1 } \gamma \Gamma)$   
**from this obtain**  $\Phi \gamma$  **where**  $\gamma: \gamma \in \text{set } \Gamma \vdash \gamma \rightarrow \sigma$   
**and**  $\Phi: \text{map snd } \Phi = (\text{remove1 } \sigma \Sigma)$   
 $\text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } (\text{remove1 } \gamma \Gamma)$   
 $\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma$   
**unfolding** *stronger-theory-relation-def* **by** *blast*  
**let**  $? \Phi = (\gamma, \sigma) \# \Phi$   
**from**  $\Phi(1)$  **have**  $\text{map snd } ? \Phi = \sigma \# (\text{remove1 } \sigma \Sigma)$  **by** *simp*  
**moreover from**  $\Phi(2) \gamma(1)$  **have**  $\text{mset } (\text{map fst } ? \Phi) \subseteq\# \text{mset } \Gamma$   
**by** (*simp add: insert-subset-eq-iff*)  
**moreover from**  $\Phi(3) \gamma(2)$  **have**  $\forall (\gamma, \sigma) \in \text{set } ? \Phi. \vdash \gamma \rightarrow \sigma$   
**by** *auto*  
**ultimately have**  $(\sigma \# (\text{remove1 } \sigma \Sigma)) \preceq \Gamma$   
**unfolding** *stronger-theory-relation-def* **by** *metis*  
**moreover from** *assms* **have**  $\sigma \# (\text{remove1 } \sigma \Sigma) \Rightarrow \Sigma$   
**by** *force*  
**ultimately show**  $\Sigma \preceq \Gamma$   
**using** *stronger-theory-left-permutation* **by** *blast*  
**qed**

lemma (in *implication-logic*) *stronger-theory-cons-witness*:

$(\sigma \# \Sigma) \preceq \Gamma = (\exists \gamma \in \text{set } \Gamma. \vdash \gamma \rightarrow \sigma \wedge \Sigma \preceq (\text{remove1 } \gamma \Gamma))$   
**proof** –  
**have**  $\sigma \in \# \text{mset } (\sigma \# \Sigma)$  **by** *simp*  
**hence**  $(\sigma \# \Sigma) \preceq \Gamma = (\exists \gamma \in \text{set } \Gamma. \vdash \gamma \rightarrow \sigma \wedge (\text{remove1 } \sigma (\sigma \# \Sigma))) \preceq (\text{remove1 } \gamma \Gamma)$   
**by** (*meson list.set-intros(1) stronger-theory-witness*)  
**thus** *?thesis* **by** *simp*  
**qed**

**lemma** (in *implication-logic*) *stronger-theory-left-cons*:

**assumes**  $(\sigma \# \Sigma) \preceq \Gamma$

**shows**  $\Sigma \preceq \Gamma$

**proof** –

**from** *assms* **obtain**  $\Phi$  **where**  $\Phi$ :

*map snd*  $\Phi = \sigma \# \Sigma$

*mset (map fst*  $\Phi) \subseteq \# \text{mset } \Gamma$

$\forall (\delta, \sigma) \in \text{set } \Phi. \vdash \delta \rightarrow \sigma$

**using** *stronger-theory-relation-def* **by** (*simp, metis*)

**let**  $?\Phi' = \text{remove1 } (\text{hd } \Phi) \Phi$

**from**  $\Phi(1)$  **have** *map snd*  $?\Phi' = \Sigma$  **by** (*induct*  $\Phi$ , *simp+*)

**moreover from**  $\Phi(2)$  **have** *mset (map fst*  $?\Phi') \subseteq \# \text{mset } \Gamma$

**by** (*metis diff-subset-eq-self*

*list-subtract.simps(1)*

*list-subtract.simps(2)*

*list-subtract-mset-homomorphism*

*map-monotonic*

*subset-mset.dual-order.trans*)

**moreover from**  $\Phi(3)$  **have**  $\forall (\delta, \sigma) \in \text{set } ?\Phi'. \vdash \delta \rightarrow \sigma$  **by** *fastforce*

**ultimately show** *?thesis* **unfolding** *stronger-theory-relation-def* **by** *blast*

**qed**

**lemma** (in *implication-logic*) *stronger-theory-right-cons*:

**assumes**  $\Sigma \preceq \Gamma$

**shows**  $\Sigma \preceq (\gamma \# \Gamma)$

**proof** –

**from** *assms* **obtain**  $\Phi$  **where**  $\Phi$ :

*map snd*  $\Phi = \Sigma$

*mset (map fst*  $\Phi) \subseteq \# \text{mset } \Gamma$

$\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma$

**unfolding** *stronger-theory-relation-def*

**by** *auto*

**hence** *mset (map fst*  $\Phi) \subseteq \# \text{mset } (\gamma \# \Gamma)$

**by** (*metis Diff-eq-empty-iff-mset*

*list-subtract.simps(2)*

*list-subtract-mset-homomorphism*

*mset-zero-iff remove1.simps(1)*)

**with**  $\Phi(1)$   $\Phi(3)$  **show** *?thesis*

**unfolding** *stronger-theory-relation-def*

**by** *auto*

qed

**lemma** (in *implication-logic*) *stronger-theory-left-right-cons*:

**assumes**  $\vdash \gamma \rightarrow \sigma$   
**and**  $\Sigma \preceq \Gamma$   
**shows**  $(\sigma \# \Sigma) \preceq (\gamma \# \Gamma)$

**proof** –

**from** *assms*(2) **obtain**  $\Phi$  **where**  $\Phi$ :

*map snd*  $\Phi = \Sigma$   
*mset (map fst*  $\Phi) \subseteq\#$  *mset*  $\Gamma$   
 $\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma$   
**unfolding** *stronger-theory-relation-def*  
**by** *auto*

**let**  $?\Phi = (\gamma, \sigma) \# \Phi$

**from** *assms*(1)  $\Phi$  **have**

*map snd*  $?\Phi = \sigma \# \Sigma$   
*mset (map fst*  $?\Phi) \subseteq\#$  *mset*  $(\gamma \# \Gamma)$   
 $\forall (\gamma, \sigma) \in \text{set } ?\Phi. \vdash \gamma \rightarrow \sigma$   
**by** *fastforce+*

**thus** *?thesis*

**unfolding** *stronger-theory-relation-def*  
**by** *metis*

qed

**lemma** (in *implication-logic*) *stronger-theory-relation-alt-def*:

$\Sigma \preceq \Gamma = (\exists \Phi. \text{mset } (\text{map snd } \Phi) = \text{mset } \Sigma \wedge$   
 $\text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } \Gamma \wedge$   
 $(\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma))$

**proof** (*induct*  $\Gamma$  *arbitrary*:  $\Sigma$ )

**case** *Nil*

**then show** *?case*

**using** *stronger-theory-empty-list-intro*  
*stronger-theory-reflexive*

**by** (*simp, blast*)

**next**

**case** (*Cons*  $\gamma$   $\Gamma$ )

**have**  $\Sigma \preceq (\gamma \# \Gamma) = (\exists \Phi. \text{mset } (\text{map snd } \Phi) = \text{mset } \Sigma \wedge$   
 $\text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } (\gamma \# \Gamma) \wedge$   
 $(\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma))$

**proof** (*rule iffI*)

**assume**  $\Sigma \preceq (\gamma \# \Gamma)$

**thus**  $\exists \Phi. \text{mset } (\text{map snd } \Phi) = \text{mset } \Sigma \wedge$   
 $\text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } (\gamma \# \Gamma) \wedge$   
 $(\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma)$

**unfolding** *stronger-theory-relation-def*

**by** *metis*

**next**

**assume**  $\exists \Phi. \text{mset } (\text{map snd } \Phi) = \text{mset } \Sigma \wedge$   
 $\text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } (\gamma \# \Gamma) \wedge$

$(\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma)$   
**from this obtain  $\Phi$  where  $\Phi$ :**  
 $\text{mset } (\text{map } \text{snd } \Phi) = \text{mset } \Sigma$   
 $\text{mset } (\text{map } \text{fst } \Phi) \subseteq \# \text{mset } (\gamma \# \Gamma)$   
 $\forall (\gamma, \sigma) \in \text{set } \Phi. \vdash \gamma \rightarrow \sigma$   
**by metis**  
**show  $\Sigma \preceq (\gamma \# \Gamma)$**   
**proof** (*cases*  $\exists \sigma. (\gamma, \sigma) \in \text{set } \Phi$ )  
**assume**  $\exists \sigma. (\gamma, \sigma) \in \text{set } \Phi$   
**from this obtain  $\sigma$  where  $\sigma: (\gamma, \sigma) \in \text{set } \Phi$  by auto**  
**let  $?\Phi = \text{remove1 } (\gamma, \sigma) \Phi$**   
**from  $\sigma$  have  $\text{mset } (\text{map } \text{snd } ?\Phi) = \text{mset } (\text{remove1 } \sigma \Sigma)$**   
**using  $\Phi(1)$  remove1-pairs-list-projections-snd by force+**  
**moreover**  
**from  $\sigma$  have  $\text{mset } (\text{map } \text{fst } ?\Phi) = \text{mset } (\text{remove1 } \gamma (\text{map } \text{fst } \Phi))$**   
**using  $\Phi(1)$  remove1-pairs-list-projections-fst by force+**  
**with  $\Phi(2)$  have  $\text{mset } (\text{map } \text{fst } ?\Phi) \subseteq \# \text{mset } \Gamma$**   
**by (*simp add: subset-eq-diff-conv*)**  
**moreover from  $\Phi(3)$  have  $\forall (\gamma, \sigma) \in \text{set } ?\Phi. \vdash \gamma \rightarrow \sigma$**   
**by fastforce**  
**ultimately have  $\text{remove1 } \sigma \Sigma \preceq \Gamma$  using *Cons* by blast**  
**from this obtain  $\Psi$  where  $\Psi$ :**  
 $\text{map } \text{snd } \Psi = \text{remove1 } \sigma \Sigma$   
 $\text{mset } (\text{map } \text{fst } \Psi) \subseteq \# \text{mset } \Gamma$   
 $\forall (\gamma, \sigma) \in \text{set } \Psi. \vdash \gamma \rightarrow \sigma$   
**unfolding stronger-theory-relation-def**  
**by blast**  
**let  $?\Psi = (\gamma, \sigma) \# \Psi$**   
**from  $\Psi$  have  $\text{map } \text{snd } ?\Psi = \sigma \# (\text{remove1 } \sigma \Sigma)$**   
 $\text{mset } (\text{map } \text{fst } ?\Psi) \subseteq \# \text{mset } (\gamma \# \Gamma)$   
**by simp+**  
**moreover from  $\Phi(3)$   $\sigma$  have  $\vdash \gamma \rightarrow \sigma$  by auto**  
**with  $\Psi(3)$  have  $\forall (\gamma, \sigma) \in \text{set } ?\Psi. \vdash \gamma \rightarrow \sigma$  by auto**  
**ultimately have  $(\sigma \# (\text{remove1 } \sigma \Sigma)) \preceq (\gamma \# \Gamma)$**   
**unfolding stronger-theory-relation-def**  
**by metis**  
**moreover**  
**have  $\sigma \in \text{set } \Sigma$**   
**by (*metis*  $\Phi(1)$   $\sigma$  *set-mset-mset set-zip-rightD zip-map-fst-snd*)**  
**hence  $\Sigma \rightleftharpoons \sigma \# (\text{remove1 } \sigma \Sigma)$**   
**by auto**  
**hence  $\Sigma \preceq (\sigma \# (\text{remove1 } \sigma \Sigma))$**   
**using stronger-theory-reflexive**  
*stronger-theory-right-permutation*  
**by blast**  
**ultimately show *?thesis***  
**using stronger-theory-transitive**  
**by blast**  
**next**



```

assume  $\nexists \sigma. (\gamma, \sigma) \in \text{set } \Phi$ 
hence  $\gamma \notin \text{set } (\text{map fst } \Phi)$  by fastforce
with  $\Phi(2)$  have  $\text{mset } (\text{map fst } \Phi) \subseteq\# \text{mset } \Gamma$ 
  by (metis diff-single-trivial
      in-multiset-in-set
      insert-DiffM2
      mset-remove1
      remove-hd
      subset-eq-diff-conv)
hence  $\Sigma \preceq \Gamma$ 
  using Cons  $\Phi(1)$   $\Phi(3)$ 
  by blast
thus ?thesis
  using stronger-theory-right-cons
  by auto
qed
qed
thus ?case by auto
qed

lemma (in implication-logic) stronger-theory-deduction-monotonic:
  assumes  $\Sigma \preceq \Gamma$ 
    and  $\Sigma \vdash \varphi$ 
  shows  $\Gamma \vdash \varphi$ 
using assms
proof (induct  $\Sigma$  arbitrary:  $\varphi$ )
  case Nil
  then show ?case
    by (simp add: list-deduction-weaken)
next
  case (Cons  $\sigma$   $\Sigma$ )
  assume  $(\sigma \# \Sigma) \preceq \Gamma$   $(\sigma \# \Sigma) \vdash \varphi$ 
  hence  $\Sigma \vdash \sigma \rightarrow \varphi$   $\Sigma \preceq \Gamma$ 
  using
    list-deduction-theorem
    stronger-theory-left-cons
  by (blast, metis)
  with Cons have  $\Gamma \vdash \sigma \rightarrow \varphi$  by blast
moreover
  have  $\sigma \in \text{set } (\sigma \# \Sigma)$  by auto
  with  $\langle \sigma \# \Sigma \rangle \preceq \Gamma$  obtain  $\gamma$  where  $\gamma \in \text{set } \Gamma \vdash \gamma \rightarrow \sigma$ 
  using stronger-theory-witness by blast
  hence  $\Gamma \vdash \sigma$ 
  using
    list-deduction-modus-ponens
    list-deduction-reflection
    list-deduction-weaken
  by blast
ultimately have  $\Gamma \vdash \varphi$ 

```

```

    using list-deduction-modus-ponens by blast
  then show ?case by blast
qed

lemma (in classical-logic) measure-msub-left-monotonic:
  assumes mset  $\Sigma \subseteq\#$  mset  $\Gamma$ 
    and  $\Sigma \text{\$}\vdash \Phi$ 
  shows  $\Gamma \text{\$}\vdash \Phi$ 
  using assms
proof (induct  $\Phi$  arbitrary:  $\Sigma \Gamma$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\varphi \Phi$ )
  from this obtain  $\Psi$  where  $\Psi$ :
    mset (map snd  $\Psi$ )  $\subseteq\#$  mset  $\Sigma$ 
    map (uncurry ( $\sqcup$ ))  $\Psi \text{\$}\vdash \varphi$ 
    map (uncurry ( $\rightarrow$ ))  $\Psi @ \Sigma \ominus$  (map snd  $\Psi$ )  $\text{\$}\vdash \Phi$ 
  using measure-deduction.simps(2) by blast
  let ? $\Psi$  = map snd  $\Psi$ 
  let ? $\Psi'$  = map (uncurry ( $\rightarrow$ ))  $\Psi$ 
  let ? $\Sigma'$  = ? $\Psi' @ (\Sigma \ominus ?\Psi)$ 
  let ? $\Gamma'$  = ? $\Psi' @ (\Gamma \ominus ?\Psi)$ 
  from  $\Psi$  have mset ? $\Psi \subseteq\#$  mset  $\Gamma$ 
    using  $\langle$ mset  $\Sigma \subseteq\#$  mset  $\Gamma\rangle$  subset-mset.trans by blast
  moreover have mset ( $\Sigma \ominus ?\Psi$ )  $\subseteq\#$  mset ( $\Gamma \ominus ?\Psi$ )
    by (metis  $\langle$ mset  $\Sigma \subseteq\#$  mset  $\Gamma\rangle$  list-subtract-monotonic)
  hence mset ? $\Sigma' \subseteq\#$  mset ? $\Gamma'$ 
    by simp
  with Cons.hyps  $\Psi(3)$  have ? $\Gamma' \text{\$}\vdash \Phi$  by blast
  ultimately have  $\Gamma \text{\$}\vdash (\varphi \# \Phi)$ 
    using  $\Psi(2)$  by fastforce
  then show ?case
    by simp
qed

```

```

lemma (in classical-logic) witness-weaker-theory:
  assumes mset (map snd  $\Sigma$ )  $\subseteq\#$  mset  $\Gamma$ 
  shows map (uncurry ( $\sqcup$ ))  $\Sigma \preceq \Gamma$ 
proof -
  have  $\forall \Gamma. \text{mset (map snd } \Sigma) \subseteq\# \text{ mset } \Gamma \longrightarrow \text{map (uncurry } (\sqcup)) \Sigma \preceq \Gamma$ 
  proof (induct  $\Sigma$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\sigma \Sigma$ )
    {
      fix  $\Gamma$ 
      assume mset (map snd ( $\sigma \# \Sigma$ ))  $\subseteq\#$  mset  $\Gamma$ 

```

```

hence  $mset (map\ snd\ \Sigma) \subseteq\# mset (remove1 (snd\ \sigma)\ \Gamma)$ 
  by (simp add: insert-subset-eq-iff)
with Cons have  $map (uncurry (\sqcup))\ \Sigma \preceq remove1 (snd\ \sigma)\ \Gamma$  by blast
moreover have  $uncurry (\sqcup) = (\lambda\ \sigma.\ fst\ \sigma \sqcup snd\ \sigma)$  by fastforce
hence  $uncurry (\sqcup)\ \sigma = fst\ \sigma \sqcup snd\ \sigma$  by simp
moreover have  $\vdash\ snd\ \sigma \rightarrow (fst\ \sigma \sqcup snd\ \sigma)$ 
  unfolding disjunction-def
  by (simp add: axiom-k)
ultimately have  $map (uncurry (\sqcup)) (\sigma\ \#\ \Sigma) \preceq (snd\ \sigma\ \#\ (remove1 (snd\ \sigma)\ \Gamma))$ 
 $\Gamma))$ 
  by (simp add: stronger-theory-left-right-cons)
moreover have  $mset (snd\ \sigma\ \#\ (remove1 (snd\ \sigma)\ \Gamma)) = mset\ \Gamma$ 
  using  $\langle mset (map\ snd (\sigma\ \#\ \Sigma)) \subseteq\# mset\ \Gamma \rangle$ 
  by (simp, meson insert-DiffM mset-subset-eq-insertD)
ultimately have  $map (uncurry (\sqcup)) (\sigma\ \#\ \Sigma) \preceq \Gamma$ 
  unfolding stronger-theory-relation-alt-def
  by simp
}
then show ?case by blast
qed
with assms show ?thesis by simp
qed

```

**lemma** (in *implication-logic*) *stronger-theory-combine*:

```

assumes  $\Phi \preceq \Delta$ 
  and  $\Psi \preceq \Gamma$ 
shows  $(\Phi\ @\ \Psi) \preceq (\Delta\ @\ \Gamma)$ 
proof –
have  $\forall\ \Phi.\ \Phi \preceq \Delta \longrightarrow (\Phi\ @\ \Psi) \preceq (\Delta\ @\ \Gamma)$ 
proof (induct  $\Delta$ )
  case Nil
  then show ?case
  using assms(2) stronger-theory-empty-list-intro by fastforce
next
case (Cons  $\delta\ \Delta$ )
  {
  fix  $\Phi$ 
  assume  $\Phi \preceq (\delta\ \#\ \Delta)$ 
  from this obtain  $\Sigma$  where  $\Sigma$ :
     $map\ snd\ \Sigma = \Phi$ 
     $mset (map\ fst\ \Sigma) \subseteq\# mset (\delta\ \#\ \Delta)$ 
     $\forall (\delta, \varphi) \in set\ \Sigma.\ \vdash\ \delta \rightarrow \varphi$ 
  unfolding stronger-theory-relation-def
  by blast
  have  $(\Phi\ @\ \Psi) \preceq ((\delta\ \#\ \Delta)\ @\ \Gamma)$ 
  proof (cases  $\exists\ \varphi.\ (\delta, \varphi) \in set\ \Sigma$ )
    assume  $\exists\ \varphi.\ (\delta, \varphi) \in set\ \Sigma$ 
    from this obtain  $\varphi$  where  $\varphi: (\delta, \varphi) \in set\ \Sigma$  by auto
    let  $\Sigma = remove1 (\delta, \varphi)\ \Sigma$ 

```

**from**  $\varphi \Sigma(1)$  **have**  $mset (map\ snd\ ?\Sigma) = mset (remove1\ \varphi\ \Phi)$   
**using** *remove1-pairs-list-projections-snd* **by** *fastforce*  
**moreover from**  $\varphi$  **have**  $mset (map\ fst\ ?\Sigma) = mset (remove1\ \delta (map\ fst$   
 $\Sigma))$   
**using** *remove1-pairs-list-projections-fst* **by** *fastforce*  
**hence**  $mset (map\ fst\ ?\Sigma) \subseteq\# mset\ \Delta$   
**using**  $\Sigma(2)$  *mset.simps(1)* *subset-eq-diff-conv* **by** *force*  
**moreover from**  $\Sigma(3)$  **have**  $\forall (\delta, \varphi) \in set\ ?\Sigma. \vdash \delta \rightarrow \varphi$  **by** *auto*  
**ultimately have**  $remove1\ \varphi\ \Phi \preceq \Delta$   
**unfolding** *stronger-theory-relation-alt-def* **by** *blast*  
**hence**  $(remove1\ \varphi\ \Phi @ \Psi) \preceq (\Delta @ \Gamma)$  **using** *Cons* **by** *auto*  
**from this obtain**  $\Omega$  **where**  $\Omega$ :  
 $map\ snd\ \Omega = (remove1\ \varphi\ \Phi) @ \Psi$   
 $mset (map\ fst\ \Omega) \subseteq\# mset (\Delta @ \Gamma)$   
 $\forall (\alpha, \beta) \in set\ \Omega. \vdash \alpha \rightarrow \beta$   
**unfolding** *stronger-theory-relation-def*  
**by** *blast*  
**let**  $?\Omega = (\delta, \varphi) \# \Omega$   
**have**  $map\ snd\ ?\Omega = \varphi \# remove1\ \varphi\ \Phi @ \Psi$   
**using**  $\Omega(1)$  **by** *simp*  
**moreover have**  $mset (map\ fst\ ?\Omega) \subseteq\# mset ((\delta \# \Delta) @ \Gamma)$   
**using**  $\Omega(2)$  **by** *simp*  
**moreover have**  $\vdash \delta \rightarrow \varphi$   
**using**  $\Sigma(3)$   $\varphi$  **by** *blast*  
**hence**  $\forall (\alpha, \beta) \in set\ ?\Omega. \vdash \alpha \rightarrow \beta$  **using**  $\Omega(3)$  **by** *auto*  
**ultimately have**  $(\varphi \# remove1\ \varphi\ \Phi @ \Psi) \preceq ((\delta \# \Delta) @ \Gamma)$   
**by** (*metis stronger-theory-relation-def*)  
**moreover have**  $\varphi \in set\ \Phi$   
**using**  $\Sigma(1)$   $\varphi$  **by** *force*  
**hence**  $(\varphi \# remove1\ \varphi\ \Phi) \rightleftharpoons \Phi$   
**by** *force*  
**hence**  $(\varphi \# remove1\ \varphi\ \Phi @ \Psi) \rightleftharpoons \Phi @ \Psi$   
**by** (*metis append-Cons perm-append2*)  
**ultimately show** *?thesis*  
**using** *stronger-theory-left-permutation* **by** *blast*  
**next**  
**assume**  $\nexists \varphi. (\delta, \varphi) \in set\ \Sigma$   
**hence**  $\delta \notin set (map\ fst\ \Sigma)$   
 $mset\ \Delta + add\ mset\ \delta (mset\ []) = mset (\delta \# \Delta)$   
**by** *auto*  
**hence**  $mset (map\ fst\ \Sigma) \subseteq\# mset\ \Delta$   
**by** (*metis (no-types) <mset (map fst Σ) ⊆# mset (δ # Δ)>*  
*diff-single-trivial*  
*mset.simps(1)*  
*set-mset-mset*  
*subset-eq-diff-conv*)  
**with**  $\Sigma(1)$   $\Sigma(3)$  **have**  $\Phi \preceq \Delta$   
**unfolding** *stronger-theory-relation-def*  
**by** *blast*

```

    hence  $(\Phi @ \Psi) \preceq (\Delta @ \Gamma)$  using Cons by auto
    then show ?thesis
      by (simp add: stronger-theory-right-cons)
    qed
  }
  then show ?case by blast
  qed
  thus ?thesis using assms by blast
  qed

```

We now turn to proving that  $(\succeq)$  is a subrelation of  $(:\vdash)$ .

**lemma** (in *classical-logic*) *stronger-theory-to-measure-deduction*:

```

  assumes  $\Gamma \succeq \Sigma$ 
  shows  $\Gamma \text{ \$}\vdash \Sigma$ 
  proof -
    have  $\forall \Gamma. \Sigma \preceq \Gamma \longrightarrow \Gamma \text{ \$}\vdash \Sigma$ 
    proof (induct  $\Sigma$ )
      case Nil
      then show ?case by fastforce
    next
      case (Cons  $\sigma \Sigma$ )
      {
        fix  $\Gamma$ 
        assume  $(\sigma \# \Sigma) \preceq \Gamma$ 
        from this obtain  $\gamma$  where  $\gamma: \gamma \in \text{set } \Gamma \vdash \gamma \rightarrow \sigma \Sigma \preceq (\text{remove1 } \gamma \Gamma)$ 
          using stronger-theory-cons-witness by blast
        let  $?\Phi = [(\gamma, \gamma)]$ 
        from  $\gamma$  Cons have  $(\text{remove1 } \gamma \Gamma) \text{ \$}\vdash \Sigma$  by blast
        moreover have  $\text{mset } (\text{remove1 } \gamma \Gamma) \subseteq\# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \text{ ?}\Phi @ \Gamma$ 
           $\ominus (\text{map } \text{snd } \text{ ?}\Phi))$ 
          by simp
        ultimately have  $\text{map } (\text{uncurry } (\rightarrow)) \text{ ?}\Phi @ \Gamma \ominus (\text{map } \text{snd } \text{ ?}\Phi) \text{ \$}\vdash \Sigma$ 
          using measure-msub-left-monotonic by blast
        moreover have  $\text{map } (\text{uncurry } (\sqcup)) \text{ ?}\Phi \text{ :\vdash } \sigma$ 
          by (simp, metis  $\gamma(2)$ 
            Peirces-law
            disjunction-def
            list-deduction-def
            list-deduction-modus-ponens
            list-deduction-weaken
            list-implication.simps(1)
            list-implication.simps(2))
        moreover from  $\gamma(1)$  have  $\text{mset } (\text{map } \text{snd } \text{ ?}\Phi) \subseteq\# \text{mset } \Gamma$  by simp
        ultimately have  $\Gamma \text{ \$}\vdash (\sigma \# \Sigma)$ 
          using measure-deduction.simps(2) by blast
      }
    then show ?case by blast
  qed
  thus ?thesis using assms by blast

```

qed

## 2.5 Measure Deduction is a Preorder

We next show that measure deduction is a preorder.

Reflexivity follows immediately because  $(\preceq)$  is a subrelation and is itself reflexive.

**theorem** (in *classical-logic*) *measure-reflexive*:  $\Gamma \text{ \$}\vdash \Gamma$   
**by** (*simp add: stronger-theory-to-measure-deduction*)

Transitivity is complicated. It requires constructing many witnesses and involves a lot of metatheorems. Below we provide various witness constructions that allow us to establish  $\llbracket \Gamma \text{ \$}\vdash \Lambda; \Lambda \text{ \$}\vdash \Delta \rrbracket \implies \Gamma \text{ \$}\vdash \Delta$ .

**primrec** (in *implication-logic*)

*first-component* ::  $('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} (\mathfrak{A})$

**where**

$\mathfrak{A} \Psi [] = []$

|  $\mathfrak{A} \Psi (\delta \# \Delta) =$

(*case find*  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi$  of

*None*  $\Rightarrow \mathfrak{A} \Psi \Delta$

| *Some*  $\psi \Rightarrow \psi \# (\mathfrak{A} (\text{remove1 } \psi \Psi) \Delta))$ )

**primrec** (in *implication-logic*)

*second-component* ::  $('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} (\mathfrak{B})$

**where**

$\mathfrak{B} \Psi [] = []$

|  $\mathfrak{B} \Psi (\delta \# \Delta) =$

(*case find*  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi$  of

*None*  $\Rightarrow \mathfrak{B} \Psi \Delta$

| *Some*  $\psi \Rightarrow \delta \# (\mathfrak{B} (\text{remove1 } \psi \Psi) \Delta))$ )

**lemma** (in *implication-logic*) *first-component-second-component-mset-connection*:

$mset (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{A} \Psi \Delta)) = mset (\text{map } \text{snd } (\mathfrak{B} \Psi \Delta))$

**proof** –

**have**  $\forall \Psi. mset (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{A} \Psi \Delta)) = mset (\text{map } \text{snd } (\mathfrak{B} \Psi \Delta))$

**proof** (*induct*  $\Delta$ )

**case** *Nil*

**then show** *?case by simp*

**next**

**case** (*Cons*  $\delta \Delta$ )

{

**fix**  $\Psi$

**have**  $mset (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{A} \Psi (\delta \# \Delta))) =$

$mset (\text{map } \text{snd } (\mathfrak{B} \Psi (\delta \# \Delta)))$

**proof** (*cases find*  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{None}$ )

**case** *True*

**then show** *?thesis using Cons by simp*

```

next
  case False
  from this obtain  $\psi$  where
    find  $(\lambda\psi. \text{uncurry } (\rightarrow) \psi = \text{snd } \delta) \Psi = \text{Some } \psi$ 
    uncurry  $(\rightarrow) \psi = \text{snd } \delta$ 
    using find-Some-predicate
    by fastforce
  then show ?thesis using Cons by simp
qed
}
then show ?case by blast
qed
thus ?thesis by blast
qed

lemma (in implication-logic) second-component-right-empty [simp]:
   $\mathfrak{B} \ \Delta = []$ 
  by (induct  $\Delta$ , simp+)

lemma (in implication-logic) first-component-msub:
   $\text{mset } (\mathfrak{A} \ \Psi \ \Delta) \subseteq\# \text{mset } \Psi$ 
proof -
  have  $\forall \Psi. \text{mset } (\mathfrak{A} \ \Psi \ \Delta) \subseteq\# \text{mset } \Psi$ 
  proof (induct  $\Delta$ )
    case Nil
    then show ?case by simp
  next
  case (Cons  $\delta \ \Delta$ )
  {
    fix  $\Psi$ 
    have  $\text{mset } (\mathfrak{A} \ \Psi \ (\delta \ \# \ \Delta)) \subseteq\# \text{mset } \Psi$ 
    proof (cases find  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{None}$ )
      case True
      then show ?thesis using Cons by simp
    next
    case False
    from this obtain  $\psi$  where
       $\psi: \text{find } (\lambda\psi. \text{uncurry } (\rightarrow) \psi = \text{snd } \delta) \Psi = \text{Some } \psi$ 
       $\psi \in \text{set } \Psi$ 
      using find-Some-set-membership
      by fastforce
    have  $\text{mset } (\mathfrak{A} \ (\text{remove1 } \psi \ \Psi) \ \Delta) \subseteq\# \text{mset } (\text{remove1 } \psi \ \Psi)$ 
    using Cons by metis
    thus ?thesis using  $\psi$  by (simp add: insert-subset-eq-iff)
  }
  qed
}
then show ?case by blast
qed
thus ?thesis by blast

```

qed

**lemma** (in *implication-logic*) *second-component-msub*:

$mset (\mathfrak{B} \Psi \Delta) \subseteq\# mset \Delta$

**proof** –

**have**  $\forall \Psi. mset (\mathfrak{B} \Psi \Delta) \subseteq\# mset \Delta$

**proof** (*induct*  $\Delta$ )

**case** *Nil*

**then show** *?case by simp*

**next**

**case** (*Cons*  $\delta \Delta$ )

{

**fix**  $\Psi$

**have**  $mset (\mathfrak{B} \Psi (\delta \# \Delta)) \subseteq\# mset (\delta \# \Delta)$

**using** *Cons*

**by** (*cases find*  $(\lambda \psi. (uncurry (\rightarrow)) \psi = snd \delta) \Psi = None,$

*simp,*

*metis add-mset-remove-trivial*

*diff-subset-eq-self*

*subset-mset.order-trans,*

*auto*)

}

**thus** *?case by blast*

qed

**thus** *?thesis by blast*

qed

**lemma** (in *implication-logic*) *second-component-snd-projection-msub*:

$mset (map\ snd (\mathfrak{B} \Psi \Delta)) \subseteq\# mset (map (uncurry (\rightarrow)) \Psi)$

**proof** –

**have**  $\forall \Psi. mset (map\ snd (\mathfrak{B} \Psi \Delta)) \subseteq\# mset (map (uncurry (\rightarrow)) \Psi)$

**proof** (*induct*  $\Delta$ )

**case** *Nil*

**then show** *?case by simp*

**next**

**case** (*Cons*  $\delta \Delta$ )

{

**fix**  $\Psi$

**have**  $mset (map\ snd (\mathfrak{B} \Psi (\delta \# \Delta))) \subseteq\# mset (map (uncurry (\rightarrow)) \Psi)$

**proof** (*cases find*  $(\lambda \psi. (uncurry (\rightarrow)) \psi = snd \delta) \Psi = None$ )

**case** *True*

**then show** *?thesis*

**using** *Cons by simp*

**next**

**case** *False*

**from this obtain**  $\psi$  **where**  $\psi$ :

*find*  $(\lambda \psi. (uncurry (\rightarrow)) \psi = snd \delta) \Psi = Some \psi$

**by** *auto*

**hence**  $\mathfrak{B} \Psi (\delta \# \Delta) = \delta \# (\mathfrak{B} (remove1 \psi \Psi) \Delta)$



```

    using  $\psi$  by fastforce
  with Cons have mset (map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ ))  $\subseteq\#$ 
    mset ((snd  $\delta$ ) # map (uncurry ( $\rightarrow$ )) (remove1  $\psi \Psi$ ))
    by (simp, metis mset-map mset-remove1)
  moreover from  $\psi$  have snd  $\delta =$  (uncurry ( $\rightarrow$ ))  $\psi$ 
    using find-Some-predicate by fastforce
  ultimately have
    mset (map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ ))  $\subseteq\#$ 
      mset (map (uncurry ( $\rightarrow$ )) ( $\psi \#$  (remove1  $\psi \Psi$ )))
    by simp
  thus ?thesis
    by (metis
        first-component-msub
        first-component-second-component-mset-connection
        map-monotonic)
  qed
}
thus ?case by blast
qed
thus ?thesis by blast
qed

lemma (in implication-logic) second-component-diff-msub:
  assumes mset (map snd  $\Delta$ )  $\subseteq\#$  mset (map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus$  (map snd  $\Psi$ ))
  shows mset (map snd ( $\Delta \ominus$  ( $\mathfrak{B} \Psi \Delta$ )))  $\subseteq\#$  mset ( $\Gamma \ominus$  (map snd  $\Psi$ ))
proof -
  have  $\forall \Psi \Gamma. \text{mset (map snd } \Delta) \subseteq\# \text{mset (map (uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{(map snd } \Psi))} \rightarrow$ 
    mset (map snd ( $\Delta \ominus$  ( $\mathfrak{B} \Psi \Delta$ )))  $\subseteq\#$  mset ( $\Gamma \ominus$  (map snd  $\Psi$ ))
  proof (induct  $\Delta$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\delta \Delta$ )
    {
      fix  $\Psi \Gamma$ 
      assume  $\diamond: \text{mset (map snd } (\delta \# \Delta)) \subseteq\# \text{mset (map (uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{(map snd } \Psi))}$ 
      have mset (map snd (( $\delta \# \Delta$ )  $\ominus$   $\mathfrak{B} \Psi (\delta \# \Delta)$ ))  $\subseteq\#$  mset ( $\Gamma \ominus$  map snd  $\Psi$ )
      proof (cases find ( $\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta$ )  $\Psi = \text{None}$ )
        case True
        hence  $A: \text{snd } \delta \notin \text{set (map (uncurry } (\rightarrow)) \Psi)$ 
        proof (induct  $\Psi$ )
          case Nil
          then show ?case by simp
        next
          case (Cons  $\psi \Psi$ )
          then show ?case
        end
      end
    }
  end

```

by (cases uncurry ( $\rightarrow$ )  $\psi = \text{snd } \delta$ , simp+)

qed

moreover have

$mset (\text{map } \text{snd } \Delta)$   
 $\subseteq\# mset (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map } \text{snd } \Psi) - \{\#\text{snd } \delta\# \}$

using  $\diamond$  insert-subset-eq-iff by fastforce

ultimately have

$mset (\text{map } \text{snd } \Delta)$   
 $\subseteq\# mset (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ (\text{remove1 } (\text{snd } \delta) \Gamma) \ominus \text{map } \text{snd } \Psi)$

by (metis (no-types) mset-remove1 union-code list-subtract.simps(2) list-subtract-remove1-cons-perm remove1-append)

hence  $B$ :  $mset (\text{map } \text{snd } (\Delta \ominus (\mathfrak{B} \Psi \Delta))) \subseteq\# mset (\text{remove1 } (\text{snd } \delta) \Gamma \ominus (\text{map } \text{snd } \Psi))$

using Cons by blast

have  $C$ :  $\text{snd } \delta \in\# mset (\text{snd } \delta \# \text{map } \text{snd } \Delta @ (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map } \text{snd } \Psi) \ominus (\text{snd } \delta \# \text{map } \text{snd } \Delta))$

by (meson in-multiset-in-set list.set-intros(1))

have  $mset (\text{map } \text{snd } (\delta \# \Delta))$   
 $+ (mset (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map } \text{snd } \Psi) - mset (\text{map } \text{snd } (\delta \# \Delta)))$   
 $= mset (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map } \text{snd } \Psi)$

using  $\diamond$  subset-mset.add-diff-inverse by blast

then have  $\text{snd } \delta \in\# mset (\text{map } (\text{uncurry } (\rightarrow)) \Psi) + (mset \Gamma - mset (\text{map } \text{snd } \Psi))$

using  $C$  by simp

with  $A$  have  $\text{snd } \delta \in \text{set } \Gamma$

by (metis (no-types) diff-subset-eq-self in-multiset-in-set subset-mset.add-diff-inverse union-iff)

have  $D$ :  $\mathfrak{B} \Psi \Delta = \mathfrak{B} \Psi (\delta \# \Delta)$

using  $\langle \text{find } (\lambda\psi. \text{uncurry } (\rightarrow) \psi = \text{snd } \delta) \Psi = \text{None} \rangle$

by simp

obtain  $\text{diff} :: 'a \text{ list} \Rightarrow 'a \text{ list} \Rightarrow 'a \text{ list}$  where  
 $\forall x0 x1. (\exists v2. x1 @ v2 \Rightarrow x0) = (x1 @ \text{diff } x0 x1 \Rightarrow x0)$

by moura

then have  $E$ :

$mset (\text{map } \text{snd } (\mathfrak{B} \Psi (\delta \# \Delta)))$   
 $@ \text{diff } (\text{map } (\text{uncurry } (\rightarrow)) \Psi) (\text{map } \text{snd } (\mathfrak{B} \Psi (\delta \# \Delta)))$   
 $= mset (\text{map } (\text{uncurry } (\rightarrow)) \Psi)$

by (meson second-component-snd-projection-msub mset-le-perm-append)

have  $F$ :  $\forall a m \text{ ma}. (\text{add-mset } (a::'a) m \subseteq\# \text{ma}) = (a \in\# \text{ma} \wedge m \subseteq\# \text{ma} - \{\#a\# \})$

```

    using insert-subset-eq-iff by blast
  then have snd  $\delta \in \#$  mset (map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ )
    @ diff (map (uncurry ( $\rightarrow$ ))  $\Psi$ ) (map snd ( $\mathfrak{B} \Psi (\delta \#$ 
 $\Delta$ ))))
    + mset ( $\Gamma \ominus$  map snd  $\Psi$ )
  using E  $\diamond$  by force
  then have snd  $\delta \in \#$  mset ( $\Gamma \ominus$  map snd  $\Psi$ )
  using A E by (metis (no-types) in-multiset-in-set union-iff)
  then have G: add-mset (snd  $\delta$ ) (mset (map snd ( $\Delta \ominus \mathfrak{B} \Psi \Delta$ )))  $\subseteq \#$  mset
( $\Gamma \ominus$  map snd  $\Psi$ )
  using B F by force
  have H:  $\forall ps$  psa f.  $\neg$  mset (ps::('a  $\times$  'a) list)  $\subseteq \#$  mset psa  $\vee$ 
    mset ((map f psa::'a list)  $\ominus$  map f ps) = mset (map f (psa
 $\ominus$  ps))
  using map-list-subtract-mset-equivalence by blast
  have snd  $\delta \notin \#$  mset (map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ ))
    + mset (diff (map (uncurry ( $\rightarrow$ ))  $\Psi$ ) (map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ )))
  using A E by auto
  then have add-mset (snd  $\delta$ ) (mset (map snd ( $\Delta \ominus \mathfrak{B} \Psi \Delta$ )))
    = mset (map snd ( $\delta \# \Delta$ )  $\ominus$  map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ ))
  using D H second-component-msub by auto
  then show ?thesis
  using G H by (metis (no-types) second-component-msub)
next
case False
from this obtain  $\psi$  where  $\psi$ : find ( $\lambda\psi$ . uncurry ( $\rightarrow$ )  $\psi$  = snd  $\delta$ )  $\Psi$  = Some
 $\psi$ 
  by auto
  let ? $\Psi'$  = remove1  $\psi$   $\Psi$ 
  let ? $\Gamma'$  = remove1 (snd  $\psi$ )  $\Gamma$ 
  have snd  $\delta$  = uncurry ( $\rightarrow$ )  $\psi$ 
     $\psi \in$  set  $\Psi$ 
    mset (( $\delta \# \Delta$ )  $\ominus \mathfrak{B} \Psi (\delta \# \Delta)$ ) =
    mset ( $\Delta \ominus \mathfrak{B} ?\Psi' \Delta$ )
  using  $\psi$  find-Some-predicate find-Some-set-membership
  by fastforce+
  moreover
  have mset ( $\Gamma \ominus$  map snd  $\Psi$ ) = mset (? $\Gamma' \ominus$  map snd ? $\Psi'$ )
    by (simp, metis  $\langle \psi \in$  set  $\Psi \rangle$  image-mset-add-mset in-multiset-in-set
insert-DiffM)
  moreover
  obtain search :: ('a  $\times$  'a) list  $\Rightarrow$  ('a  $\times$  'a  $\Rightarrow$  bool)  $\Rightarrow$  'a  $\times$  'a where
 $\forall xs$  P. ( $\exists x$ .  $x \in$  set xs  $\wedge$  P x) = (search xs P  $\in$  set xs  $\wedge$  P (search xs P))
  by moura
  then have  $\forall p$  ps. (find p ps  $\neq$  None  $\vee$  ( $\forall pa$ . pa  $\notin$  set ps  $\vee$   $\neg$  p pa))
     $\wedge$  (find p ps = None  $\vee$  search ps p  $\in$  set ps  $\wedge$  p (search ps p))
  by (metis (full-types) find-None-iff)
  then have (find ( $\lambda p$ . uncurry ( $\rightarrow$ ) p = snd  $\delta$ )  $\Psi \neq$  None
 $\vee$  ( $\forall p$ . p  $\notin$  set  $\Psi \vee$  uncurry ( $\rightarrow$ ) p  $\neq$  snd  $\delta$ ))

```

$\wedge (\text{find } (\lambda p. \text{uncurry } (\rightarrow) p = \text{snd } \delta) \Psi = \text{None}$   
 $\vee \text{search } \Psi (\lambda p. \text{uncurry } (\rightarrow) p = \text{snd } \delta) \in \text{set } \Psi$   
 $\wedge \text{uncurry } (\rightarrow) (\text{search } \Psi (\lambda p. \text{uncurry } (\rightarrow) p = \text{snd } \delta)) = \text{snd } \delta)$

**by** *blast*  
**hence**  $\text{snd } \delta \in \text{set } (\text{map } (\text{uncurry } (\rightarrow)) \Psi)$   
**by** (*metis* (*no-types*) *False image-eqI image-set*)  
**moreover**  
**have**  $A: \text{add-mset } (\text{uncurry } (\rightarrow) \psi) (\text{image-mset } \text{snd } (\text{mset } \Delta))$   
 $= \text{image-mset } \text{snd } (\text{add-mset } \delta (\text{mset } \Delta))$   
**by** (*simp add:*  $\langle \text{snd } \delta = \text{uncurry } (\rightarrow) \psi \rangle$ )  
**have**  $B: \{\#\text{snd } \delta\# \} \subseteq \# \text{image-mset } (\text{uncurry } (\rightarrow)) (\text{mset } \Psi)$   
**using**  $\langle \text{snd } \delta \in \text{set } (\text{map } (\text{uncurry } (\rightarrow)) \Psi) \rangle$  **by** *force*  
**have**  $\text{image-mset } (\text{uncurry } (\rightarrow)) (\text{mset } \Psi) - \{\#\text{snd } \delta\# \}$   
 $= \text{image-mset } (\text{uncurry } (\rightarrow)) (\text{mset } (\text{remove1 } \psi \Psi))$   
**by** (*simp add:*  $\langle \psi \in \text{set } \Psi \rangle \langle \text{snd } \delta = \text{uncurry } (\rightarrow) \psi \rangle$  *image-mset-Diff*)  
**then have**  $\text{mset } (\text{map } \text{snd } (\Delta \ominus \mathfrak{B} (\text{remove1 } \psi \Psi) \Delta))$   
 $\subseteq \# \text{mset } (\text{remove1 } (\text{snd } \psi) \Gamma \ominus \text{map } \text{snd } (\text{remove1 } \psi \Psi))$   
**by** (*metis* (*no-types*)  
 $A B \diamond \text{Cons.hyps}$   
 $\text{calculation}(1)$   
 $\text{calculation}(4)$   
 $\text{insert-subset-eq-iff}$   
 $\text{mset.simps}(2)$   
 $\text{mset-map}$   
 $\text{subset-mset.diff-add-assoc2}$   
 $\text{union-code}$ )  
**ultimately show** *?thesis* **by** *fastforce*  
**qed**  
**}**  
**then show** *?case* **by** *blast*  
**qed**  
**thus** *?thesis* **using** *assms* **by** *auto*  
**qed**

**primrec** (*in classical-logic*)  
 $\text{merge-witness} :: ('a \times 'a) \text{list} \Rightarrow ('a \times 'a) \text{list} \Rightarrow ('a \times 'a) \text{list} (\langle \mathfrak{J} \rangle)$   
**where**  
 $\mathfrak{J} \Psi [] = \Psi$   
 $|\ \mathfrak{J} \Psi (\delta \# \Delta) =$   
 $(\text{case } \text{find } (\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi \text{ of}$   
 $\text{None} \Rightarrow \delta \# \mathfrak{J} \Psi \Delta$   
 $|\ \text{Some } \psi \Rightarrow (\text{fst } \delta \sqcap \text{fst } \psi, \text{snd } \psi) \# (\mathfrak{J} (\text{remove1 } \psi \Psi) \Delta))$

**lemma** (*in classical-logic*) *merge-witness-right-empty* [*simp*]:  
 $\mathfrak{J} [] \Delta = \Delta$   
**by** (*induct*  $\Delta$ , *simp+*)

**lemma** (*in classical-logic*) *second-component-merge-witness-snd-projection*:  
 $\text{mset } (\text{map } \text{snd } \Psi @ \text{map } \text{snd } (\Delta \ominus (\mathfrak{B} \Psi \Delta))) = \text{mset } (\text{map } \text{snd } (\mathfrak{J} \Psi \Delta))$

```

proof –
  have  $\forall \Psi. \text{mset} (\text{map snd } \Psi @ \text{map snd } (\Delta \ominus (\mathfrak{B} \Psi \Delta))) = \text{mset} (\text{map snd } (\mathfrak{J} \Psi \Delta))$ 
proof (induct  $\Delta$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\delta \Delta$ )
  {
    fix  $\Psi$ 
    have  $\text{mset} (\text{map snd } \Psi @ \text{map snd } ((\delta \# \Delta) \ominus \mathfrak{B} \Psi (\delta \# \Delta))) = \text{mset} (\text{map snd } (\mathfrak{J} \Psi (\delta \# \Delta)))$ 
    proof (cases find  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{None}$ )
    case True
    then show ?thesis
    using Cons
    by (simp,
      metis (no-types, lifting)
      ab-semigroup-add-class.add-ac(1)
      add-mset-add-single
      image-mset-single
      image-mset-union
      second-component-msub
      subset-mset.add-diff-assoc2)
    next
    case False
    from this obtain  $\psi$  where  $\psi: \text{find } (\lambda \psi. \text{uncurry } (\rightarrow) \psi = \text{snd } \delta) \Psi = \text{Some } \psi$ 
    by auto
    moreover have  $\psi \in \text{set } \Psi$ 
    by (meson  $\psi$  find-Some-set-membership)
    moreover
    let  $?\Psi' = \text{remove1 } \psi \Psi$ 
    from Cons have
       $\text{mset} (\text{map snd } ?\Psi' @ \text{map snd } (\Delta \ominus \mathfrak{B} ?\Psi' \Delta)) = \text{mset} (\text{map snd } (\mathfrak{J} ?\Psi' \Delta))$ 
    by blast
    ultimately show ?thesis
    by (simp,
      metis (no-types, lifting)
      add-mset-remove-trivial-eq
      image-mset-add-mset
      in-multiset-in-set
      union-mset-add-mset-left)
    qed
  }
  then show ?case by blast
qed
thus ?thesis by blast

```

qed

**lemma** (in *classical-logic*) *second-component-merge-witness-stronger-theory*:  
 $(\text{map } (\text{uncurry } (\rightarrow)) \Delta @ \text{map } (\text{uncurry } (\rightarrow)) \Psi \ominus \text{map } \text{snd } (\mathfrak{B} \Psi \Delta)) \preceq$   
 $\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{J} \Psi \Delta)$

**proof** –

**have**  $\forall \Psi. (\text{map } (\text{uncurry } (\rightarrow)) \Delta @$   
 $\text{map } (\text{uncurry } (\rightarrow)) \Psi \ominus \text{map } \text{snd } (\mathfrak{B} \Psi \Delta)) \preceq$   
 $\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{J} \Psi \Delta)$

**proof** (*induct*  $\Delta$ )

**case** *Nil*

**then show** *?case*

**by** *simp*

**next**

**case** (*Cons*  $\delta \Delta$ )

{

**fix**  $\Psi$

**have**  $\vdash (\text{uncurry } (\rightarrow)) \delta \rightarrow (\text{uncurry } (\rightarrow)) \delta$

**using** *axiom-k modus-ponens implication-absorption* **by** *blast*

**have**

$(\text{map } (\text{uncurry } (\rightarrow)) (\delta \# \Delta) @$   
 $\text{map } (\text{uncurry } (\rightarrow)) \Psi \ominus \text{map } \text{snd } (\mathfrak{B} \Psi (\delta \# \Delta))) \preceq$   
 $\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{J} \Psi (\delta \# \Delta))$

**proof** (*cases find*  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{None}$ )

**case** *True*

**thus** *?thesis*

**using** *Cons*

$\langle \vdash (\text{uncurry } (\rightarrow)) \delta \rightarrow (\text{uncurry } (\rightarrow)) \delta \rangle$

**by** (*simp, metis stronger-theory-left-right-cons*)

**next**

**case** *False*

**from this obtain**  $\psi$  **where**  $\psi: \text{find } (\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{Some}$

$\psi$

**by** *auto*

**from**  $\psi$  **have**  $\text{snd } \delta = \text{uncurry } (\rightarrow) \psi$

**using** *find-Some-predicate* **by** *fastforce*

**from**  $\psi \langle \text{snd } \delta = \text{uncurry } (\rightarrow) \psi \rangle$  **have**

$\text{mset } (\text{map } (\text{uncurry } (\rightarrow)) (\delta \# \Delta) @$   
 $\text{map } (\text{uncurry } (\rightarrow)) \Psi \ominus \text{map } \text{snd } (\mathfrak{B} \Psi (\delta \# \Delta))) =$   
 $\text{mset } (\text{map } (\text{uncurry } (\rightarrow)) (\delta \# \Delta) @$   
 $\text{map } (\text{uncurry } (\rightarrow)) (\text{remove1 } \psi \Psi) \ominus$   
 $\text{map } \text{snd } (\mathfrak{B} (\text{remove1 } \psi \Psi) \Delta))$

**by** (*simp add: find-Some-set-membership image-mset-Diff*)

**hence**

$(\text{map } (\text{uncurry } (\rightarrow)) (\delta \# \Delta) @$   
 $\text{map } (\text{uncurry } (\rightarrow)) \Psi \ominus \text{map } \text{snd } (\mathfrak{B} \Psi (\delta \# \Delta))) \preceq$   
 $(\text{map } (\text{uncurry } (\rightarrow)) (\delta \# \Delta) @$   
 $\text{map } (\text{uncurry } (\rightarrow)) (\text{remove1 } \psi \Psi) \ominus \text{map } \text{snd } (\mathfrak{B} (\text{remove1 } \psi \Psi) \Delta))$   
**by** (*simp add: msub-stronger-theory-intro*)

```

with Cons  $\vdash$  (uncurry ( $\rightarrow$ ))  $\delta \rightarrow$  (uncurry ( $\rightarrow$ ))  $\delta$  have
  (map (uncurry ( $\rightarrow$ )) ( $\delta \# \Delta$ ) @
    map (uncurry ( $\rightarrow$ ))  $\Psi \ominus$  map snd ( $\mathfrak{B} \Psi (\delta \# \Delta)$ ))
   $\preceq$  ((uncurry ( $\rightarrow$ ))  $\delta \#$  map (uncurry ( $\rightarrow$ )) ( $\mathfrak{J}$  (remove1  $\psi \Psi$ )  $\Delta$ ))
using stronger-theory-left-right-cons
      stronger-theory-transitive
by fastforce
moreover
let  $?\alpha =$  fst  $\delta$ 
let  $?\beta =$  fst  $\psi$ 
let  $?\gamma =$  snd  $\psi$ 
have uncurry ( $\rightarrow$ ) = ( $\lambda \delta. \text{fst } \delta \rightarrow \text{snd } \delta$ ) by fastforce
with  $\psi$  have (uncurry ( $\rightarrow$ ))  $\delta = ?\alpha \rightarrow ?\beta \rightarrow ?\gamma$ 
  using find-Some-predicate by fastforce
hence  $\vdash$  (( $?\alpha \sqcap ?\beta$ )  $\rightarrow ?\gamma$ )  $\rightarrow$  (uncurry ( $\rightarrow$ ))  $\delta$ 
  using biconditional-def curry-uncurry by auto
with  $\psi$  have
  ((uncurry ( $\rightarrow$ ))  $\delta \#$  map (uncurry ( $\rightarrow$ )) ( $\mathfrak{J}$  (remove1  $\psi \Psi$ )  $\Delta$ ))  $\preceq$ 
  map (uncurry ( $\rightarrow$ )) ( $\mathfrak{J} \Psi (\delta \# \Delta)$ )
  using stronger-theory-left-right-cons by auto
ultimately show  $?thesis$ 
  using stronger-theory-transitive
  by blast
qed
}
then show  $?case$  by simp
qed
thus  $?thesis$  by simp
qed

lemma (in classical-logic) merge-witness-msub-intro:
  assumes mset (map snd  $\Psi$ )  $\subseteq\#$  mset  $\Gamma$ 
  and mset (map snd  $\Delta$ )  $\subseteq\#$  mset (map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus$  (map snd
 $\Psi$ ))
  shows mset (map snd ( $\mathfrak{J} \Psi \Delta$ ))  $\subseteq\#$  mset  $\Gamma$ 
proof –
  have  $\forall \Psi \Gamma. \text{mset } (\text{map } \text{snd } \Psi) \subseteq\# \text{mset } \Gamma \longrightarrow$ 
    mset (map snd  $\Delta$ )  $\subseteq\#$  mset (map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus$  (map snd
 $\Psi$ ))  $\longrightarrow$ 
    mset (map snd ( $\mathfrak{J} \Psi \Delta$ ))  $\subseteq\#$  mset  $\Gamma$ 
proof (induct  $\Delta$ )
  case Nil
  then show  $?case$  by simp
next
case (Cons  $\delta \Delta$ )
  {
  fix  $\Psi :: ('a \times 'a)$  list
  fix  $\Gamma :: 'a$  list
  assume  $\diamond: \text{mset } (\text{map } \text{snd } \Psi) \subseteq\# \text{mset } \Gamma$ 

```

```

      mset (map snd (δ # Δ)) ⊆# mset (map (uncurry (→)) Ψ @ Γ ⊖
(map snd Ψ))
  have mset (map snd (⋈ Ψ (δ # Δ))) ⊆# mset Γ
  proof (cases find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None)
    case True
      hence snd δ ∉ set (map (uncurry (→)) Ψ)
      proof (induct Ψ)
        case Nil
          then show ?case by simp
        next
          case (Cons ψ Ψ)
            hence uncurry (→) ψ ≠ snd δ by fastforce
            with Cons show ?case by fastforce
      qed
  with ◇(2) have snd δ ∈# mset (Γ ⊖ map snd Ψ)
    using mset-subset-eq-insertD by fastforce
  with ◇(1) have mset (map snd Ψ) ⊆# mset (remove1 (snd δ) Γ)
    by (metis list-subtract-mset-homomorphism
      mset-remove1
      single-subset-iff
      subset-mset.add-diff-assoc
      subset-mset.add-diff-inverse
      subset-mset.le-iff-add)
  moreover
  have add-mset (snd δ) (mset (Γ ⊖ map snd Ψ) - {#snd δ#}) = mset (Γ
⊖ map snd Ψ)
    by (meson ⟨snd δ ∈# mset (Γ ⊖ map snd Ψ)⟩ insert-DiffM)
    then have image-mset snd (mset Δ) - (mset Γ - add-mset (snd δ)
(image-mset snd (mset Ψ)))
      ⊆# {#x → y. (x, y) ∈# mset Ψ#}
    using ◇(2) by (simp, metis add-mset-diff-bothsides
      list-subtract-mset-homomorphism
      mset-map subset-eq-diff-conv)
  hence mset (map snd Δ)
    ⊆# mset (map (uncurry (→)) Ψ @ (remove1 (snd δ) Γ) ⊖ (map snd Ψ))
    using subset-eq-diff-conv by (simp, blast)
  ultimately have mset (map snd (⋈ Ψ Δ)) ⊆# mset (remove1 (snd δ) Γ)
    using Cons by blast
  hence mset (map snd (δ # (⋈ Ψ Δ))) ⊆# mset Γ
    by (simp, metis ⟨snd δ ∈# mset (Γ ⊖ map snd Ψ)⟩
      cancel-ab-semigroup-add-class.diff-right-commute
      diff-single-trivial
      insert-subset-eq-iff
      list-subtract-mset-homomorphism
      multi-drop-mem-not-eq)
  with ⟨find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None⟩
  show ?thesis
    by simp
  next

```



```

case False
from this obtain  $\psi$  where  $\psi$ :
  find  $(\lambda\psi. \text{uncurry } (\rightarrow) \psi = \text{snd } \delta) \Psi = \text{Some } \psi$ 
  by fastforce
let  $?x = \text{fst } \psi$ 
let  $?y = \text{snd } \psi$ 
have  $\text{uncurry } (\rightarrow) = (\lambda \psi. \text{fst } \psi \rightarrow \text{snd } \psi)$ 
  by fastforce
moreover
from this have  $\text{uncurry } (\rightarrow) \psi = ?x \rightarrow ?y$  by fastforce
with  $\psi$  have  $A: (?x, ?y) \in \text{set } \Psi$ 
  and  $B: \text{snd } \delta = ?x \rightarrow ?y$ 
  using find-Some-predicate
  by (simp add: find-Some-set-membership, fastforce)
let  $?x' = \text{remove1 } (?x, ?y) \Psi$ 
from  $B \diamond(2)$  have
   $\text{mset } (\text{map } \text{snd } \Delta) \subseteq\# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map } \text{snd } \Psi)$ 
-  $\{ \# ?x \rightarrow ?y \# \}$ 
  by (simp add: insert-subset-eq-iff)
moreover
have  $\text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \Psi)$ 
  =  $\text{add-mset } (\text{case } (\text{fst } \psi, \text{snd } \psi) \text{ of } (x, xa) \Rightarrow x \rightarrow xa)$ 
  ( $\text{image-mset } (\text{uncurry } (\rightarrow)) (\text{mset } (\text{remove1 } (\text{fst } \psi, \text{snd } \psi) \Psi))$ )
  by (metis (no-types)
     $A$ 
    image-mset-add-mset
    in-multiset-in-set
    insert-DiffM
    mset-map
    mset-remove1
    uncurry-def)
ultimately have
 $\text{mset } (\text{map } \text{snd } \Delta) \subseteq\# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) ?x' @ \Gamma \ominus \text{map } \text{snd } \Psi)$ 
using
  add-diff-cancel-left'
  add-diff-cancel-right
  diff-diff-add-mset
  diff-subset-eq-self
  mset-append
  subset-eq-diff-conv
  subset-mset.diff-add
by auto
moreover from  $A B \diamond$ 
have  $\text{mset } (\Gamma \ominus \text{map } \text{snd } \Psi) = \text{mset}((\text{remove1 } ?y \Gamma) \ominus (\text{remove1 } ?y (\text{map } \text{snd } \Psi)))$ 
using
  image-eqI
  prod.sel(2)
  set-map

```

**by** *force*  
**with** *A* **have**  
 $mset (\Gamma \ominus \text{map } \text{snd } \Psi) = mset((\text{remove1 } ?\gamma \Gamma) \ominus (\text{map } \text{snd } ?\Psi'))$   
**by** (*metis*  
 $\text{remove1-pairs-list-projections-snd}$   
 $\text{in-multiset-in-set}$   
 $\text{list-subtract-mset-homomorphism}$   
 $\text{mset-remove1}$ )  
**ultimately have**  
 $mset (\text{map } \text{snd } \Delta) \subseteq\# mset (\text{map } (\text{uncurry } (\rightarrow)) ?\Psi'$   
 $\quad \text{@ } (\text{remove1 } ?\gamma \Gamma)$   
 $\quad \ominus \text{map } \text{snd } ?\Psi')$   
**by** *simp*  
**hence**  $mset (\text{map } \text{snd } (\mathfrak{J} ?\Psi' \Delta)) \subseteq\# mset (\text{remove1 } ?\gamma \Gamma)$   
**using** *Cons*  $\diamond(1)$  *A*  
**by** (*metis* (*no-types*, *lifting*)  
 $\text{image-mset-add-mset}$   
 $\text{in-multiset-in-set}$   
 $\text{insert-DiffM}$   
 $\text{insert-subset-eq-iff}$   
 $\text{mset-map mset-remove1}$   
 $\text{prod.collapse}$ )  
**with**  $\diamond(1)$  *A* **have**  $mset (\text{map } \text{snd } (\mathfrak{J} ?\Psi' \Delta)) + \{\# ?\gamma \#\} \subseteq\# mset \Gamma$   
**by** (*metis* *add-mset-add-single*  
 $\text{image-eqI}$   
 $\text{insert-subset-eq-iff}$   
 $\text{mset-remove1}$   
 $\text{mset-subset-eqD}$   
 $\text{set-map}$   
 $\text{set-mset-mset}$   
 $\text{snd-conv}$ )  
**hence**  $mset (\text{map } \text{snd } ((\text{fst } \delta \sqcap ?\chi, ?\gamma) \# (\mathfrak{J} ?\Psi' \Delta))) \subseteq\# mset \Gamma$   
**by** *simp*  
**moreover from**  $\psi$  **have**  
 $\mathfrak{J} \Psi (\delta \# \Delta) = (\text{fst } \delta \sqcap ?\chi, ?\gamma) \# (\mathfrak{J} ?\Psi' \Delta)$   
**by** *simp*  
**ultimately show** *?thesis* **by** *simp*  
**qed**  
**}**  
**thus** *?case* **by** *blast*  
**qed**  
**with** *assms* **show** *?thesis* **by** *blast*  
**qed**

**lemma** (*in classical-logic*) *right-merge-witness-stronger-theory*:  
 $\text{map } (\text{uncurry } (\sqcup)) \Delta \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{J} \Psi \Delta)$   
**proof** –  
**have**  $\forall \Psi. \text{map } (\text{uncurry } (\sqcup)) \Delta \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{J} \Psi \Delta)$   
**proof** (*induct*  $\Delta$ )

```

case Nil
then show ?case by simp
next
case (Cons δ Δ)
{
  fix Ψ
  have map (uncurry (⊔)) (δ # Δ) ≤ map (uncurry (⊔)) (⋈ Ψ (δ # Δ))
  proof (cases find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None)
    case True
    hence ⋈ Ψ (δ # Δ) = δ # ⋈ Ψ Δ
    by simp
    moreover have ⊢ (uncurry (⊔)) δ → (uncurry (⊔)) δ
    by (metis axiom-k axiom-s modus-ponens)
    ultimately show ?thesis using Cons
    by (simp add: stronger-theory-left-right-cons)
  next
  case False
  from this obtain ψ where ψ:
    find (λ ψ. uncurry (→) ψ = snd δ) Ψ = Some ψ
  by fastforce
  let ?χ = fst ψ
  let ?γ = snd ψ
  let ?μ = fst δ
  have uncurry (→) = (λ ψ. fst ψ → snd ψ)
    uncurry (⊔) = (λ δ. fst δ ⊔ snd δ)
  by fastforce+
  hence uncurry (⊔) δ = ?μ ⊔ (?χ → ?γ)
  using ψ find-Some-predicate
  by fastforce
  moreover
  {
    fix μ χ γ
    have ⊢ ((μ ⊓ χ) ⊔ γ) → (μ ⊔ (χ → γ))
    proof -
      have ∀ M. M ⊨prop (((⟨μ⟩ ⊓ ⟨χ⟩) ⊔ ⟨γ⟩) → (⟨μ⟩ ⊔ (⟨χ⟩ → ⟨γ⟩)))
      by fastforce
      hence ⊢ (⟨((μ ⊓ χ) ⊔ γ)⟩ → (⟨μ⟩ ⊔ (⟨χ⟩ → ⟨γ⟩)))
      using propositional-semantic by blast
      thus ?thesis
      by simp
    qed
  }
  ultimately show ?thesis
  using Cons ψ stronger-theory-left-right-cons
  by simp
qed
}
thus ?case by blast
qed

```

```

thus ?thesis by blast
qed

lemma (in classical-logic) left-merge-witness-stronger-theory:
  map (uncurry (⊔)) Ψ ≲ map (uncurry (⊔)) (⋈ Ψ Δ)
proof –
  have ∀ Ψ. map (uncurry (⊔)) Ψ ≲ map (uncurry (⊔)) (⋈ Ψ Δ)
  proof (induct Δ)
    case Nil
    then show ?case
    by simp
  next
  case (Cons δ Δ)
  {
    fix Ψ
    have map (uncurry (⊔)) Ψ ≲ map (uncurry (⊔)) (⋈ Ψ (δ # Δ))
    proof (cases find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None)
      case True
      then show ?thesis
      using Cons stronger-theory-right-cons
      by auto
    next
    case False
    from this obtain ψ where ψ:
      find (λ ψ. uncurry (→) ψ = snd δ) Ψ = Some ψ
    by fastforce
    let ?χ = fst ψ
    let ?γ = snd ψ
    let ?μ = fst δ
    have uncurry (→) = (λ ψ. fst ψ → snd ψ)
      uncurry (⊔) = (λ δ. fst δ ⊔ snd δ)
    by fastforce+
    hence
      uncurry (⊔) δ = ?μ ⊔ (?χ → ?γ)
      uncurry (⊔) ψ = ?χ ⊔ ?γ
    using ψ find-Some-predicate
    by fastforce+
    moreover
    {
      fix μ χ γ
      have ⊢ ((μ ⊓ χ) ⊔ γ) → (χ ⊔ γ)
      proof –
      have ∀ ℳ. ℳ ⊨prop ((⟨μ⟩ ⊓ ⟨χ⟩) ⊔ ⟨γ⟩) → (⟨χ⟩ ⊔ ⟨γ⟩)
      by fastforce
      hence ⊢ ⟨ ((μ) ⊓ ⟨χ⟩) ⊔ ⟨γ⟩ → (⟨χ⟩ ⊔ ⟨γ⟩) ⟩
      using propositional-semantic by blast
      thus ?thesis
      by simp
    }
  }
qed

```

```

}
ultimately have
  map (uncurry (⊔)) (ψ # (remove1 ψ Ψ)) ≼
  map (uncurry (⊔)) (⋈ Ψ (δ # Δ))
using Cons ψ stronger-theory-left-right-cons
by simp
moreover from ψ have ψ ∈ set Ψ
by (simp add: find-Some-set-membership)
hence mset (map (uncurry (⊔)) (ψ # (remove1 ψ Ψ))) =
  mset (map (uncurry (⊔)) Ψ)
by (metis insert-DiffM
  mset.simps(2)
  mset-map
  mset-remove1
  set-mset-mset)
hence map (uncurry (⊔)) Ψ ≼ map (uncurry (⊔)) (ψ # (remove1 ψ Ψ))
by (simp add: msub-stronger-theory-intro)
ultimately show ?thesis
using stronger-theory-transitive by blast
qed
}
then show ?case by blast
qed
thus ?thesis by blast
qed

```

**lemma** (in *classical-logic*) *measure-empty-deduction*:

```

[] $⊢ Φ = (∀ φ ∈ set Φ. ⊢ φ)
by (induct Φ, simp, rule iffI, fastforce+)

```

**lemma** (in *classical-logic*) *measure-stronger-theory-left-monotonic*:

```

assumes Σ ≼ Γ
and Σ $⊢ Φ
shows Γ $⊢ Φ
using assms
proof (induct Φ arbitrary: Σ Γ)
case Nil
then show ?case by simp
next
case (Cons φ Φ)
from this obtain Ψ Δ where
  Ψ: mset (map snd Ψ) ⊆# mset Σ
  map (uncurry (⊔)) Ψ ⊢ φ
  map (uncurry (→)) Ψ @ Σ ⊖ (map snd Ψ) $⊢ Φ
and
  Δ: map snd Δ = Σ
  mset (map fst Δ) ⊆# mset Γ
  ∀ (γ,σ) ∈ set Δ. ⊢ γ → σ
unfolding stronger-theory-relation-def

```

**by fastforce**  
**from**  $\langle \text{mset} (\text{map snd } \Psi) \subseteq\# \text{mset } \Sigma \rangle$   
 $\langle \text{map snd } \Delta = \Sigma \rangle$   
**obtain**  $\Omega$  **where**  $\Omega$ :  
 $\text{map } (\lambda (\psi, \sigma, -). (\psi, \sigma)) \Omega = \Psi$   
 $\text{mset} (\text{map } (\lambda (-, \sigma, \gamma). (\gamma, \sigma)) \Omega) \subseteq\# \text{mset } \Delta$   
**using** *triple-list-exists* **by** *blast*  
**let**  $?\Theta = \text{map } (\lambda (\psi, -, \gamma). (\psi, \gamma)) \Omega$   
**have**  $\text{map snd } ?\Theta = \text{map fst } (\text{map } (\lambda (-, \sigma, \gamma). (\gamma, \sigma)) \Omega)$   
**by** *auto*  
**hence**  $\text{mset} (\text{map snd } ?\Theta) \subseteq\# \text{mset } \Gamma$   
**using**  $\Omega(2)$   $\Delta(2)$  *map-monotonic subset-mset.order-trans*  
**by** *metis*  
**moreover** **have**  $\text{map} (\text{uncurry } (\sqcup)) \Psi \preceq \text{map} (\text{uncurry } (\sqcup)) ?\Theta$   
**proof** –  
**let**  $?\Phi = \text{map } (\lambda (\psi, \sigma, \gamma). (\psi \sqcup \gamma, \psi \sqcup \sigma)) \Omega$   
**have**  $\text{map snd } ?\Phi = \text{map} (\text{uncurry } (\sqcup)) \Psi$   
**using**  $\Omega(1)$  **by** *fastforce*  
**moreover** **have**  $\text{map fst } ?\Phi = \text{map} (\text{uncurry } (\sqcup)) ?\Theta$   
**by** *fastforce*  
**hence**  $\text{mset} (\text{map fst } ?\Phi) \subseteq\# \text{mset} (\text{map} (\text{uncurry } (\sqcup)) ?\Theta)$   
**by** (*metis subset-mset.dual-order.refl*)  
**moreover**  
**have**  $\text{mset} (\text{map } (\lambda(\psi, \sigma, -). (\psi, \sigma)) \Omega) \subseteq\# \text{mset } \Psi$   
**using**  $\Omega(1)$  **by** *simp*  
**hence**  $\forall (\varphi, \chi) \in \text{set } ?\Phi. \vdash \varphi \rightarrow \chi$  **using**  $\Omega(2)$   
**proof** (*induct*  $\Omega$ )  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons*  $\omega$   $\Omega$ )  
**let**  $?\Phi = \text{map } (\lambda (\psi, \sigma, \gamma). (\psi \sqcup \gamma, \psi \sqcup \sigma)) (\omega \# \Omega)$   
**let**  $?\Phi' = \text{map } (\lambda (\psi, \sigma, \gamma). (\psi \sqcup \gamma, \psi \sqcup \sigma)) \Omega$   
**have**  $\text{mset} (\text{map } (\lambda(\psi, \sigma, -). (\psi, \sigma)) \Omega) \subseteq\# \text{mset } \Psi$   
 $\text{mset} (\text{map } (\lambda(-, \sigma, \gamma). (\gamma, \sigma)) \Omega) \subseteq\# \text{mset } \Delta$   
**using** *Cons.premis(1)* *Cons.premis(2)* *subset-mset.dual-order.trans* **by** *fast-*  
*force+*  
**with** *Cons* **have**  $\forall (\varphi, \chi) \in \text{set } ?\Phi'. \vdash \varphi \rightarrow \chi$  **by** *fastforce*  
**moreover**  
**let**  $?\psi = (\lambda (\psi, -, -). \psi) \omega$   
**let**  $?\sigma = (\lambda (-, \sigma, -). \sigma) \omega$   
**let**  $??\gamma = (\lambda (-, -, \gamma). \gamma) \omega$   
**have**  $(\lambda(-, \sigma, \gamma). (\gamma, \sigma)) = (\lambda \omega. ((\lambda (-, -, \gamma). \gamma) \omega, (\lambda (-, \sigma, -). \sigma) \omega))$  **by** *auto*  
**hence**  $(\lambda(-, \sigma, \gamma). (\gamma, \sigma)) \omega = (??\gamma, ?\sigma)$  **by** *metis*  
**hence**  $\vdash ?\gamma \rightarrow ?\sigma$   
**using** *Cons.premis(2)* *mset-subset-eqD*  $\Delta(3)$   
**by** *fastforce*  
**hence**  $\vdash (?\psi \sqcup ?\gamma) \rightarrow (?\psi \sqcup ?\sigma)$   
**unfolding** *disjunction-def*

**using** *modus-ponens hypothetical-syllogism*  
**by** *blast*  
**moreover have**  
 $(\lambda(\psi, \sigma, \gamma). (\psi \sqcup \gamma, \psi \sqcup \sigma)) =$   
 $(\lambda \omega. (((\lambda (\psi, -, -). \psi) \omega) \sqcup ((\lambda (-, -, \gamma). \gamma) \omega),$   
 $((\lambda (\psi, -, -). \psi) \omega) \sqcup ((\lambda (-, \sigma, -). \sigma) \omega)))$   
**by** *auto*  
**hence**  $(\lambda(\psi, \sigma, \gamma). (\psi \sqcup \gamma, \psi \sqcup \sigma)) \omega = ((? \psi \sqcup ? \gamma), (? \psi \sqcup ? \sigma))$  **by** *metis*  
**ultimately show** *?case by simp*  
**qed**  
**ultimately show** *?thesis*  
**unfolding** *stronger-theory-relation-def*  
**by** *blast*  
**qed**  
**hence**  $\text{map } (\text{uncurry } (\sqcup)) \ ?\Theta \vdash \varphi$   
**using**  $\Psi(2)$   
*stronger-theory-deduction-monotonic*  
**[where**  $\Sigma = \text{map } (\text{uncurry } (\sqcup)) \ \Psi$   
**and**  $\Gamma = \text{map } (\text{uncurry } (\sqcup)) \ ?\Theta$   
**and**  $\varphi = \varphi]$   
**by** *metis*  
**moreover have**  
 $(\text{map } (\text{uncurry } (\rightarrow)) \ \Psi @ \Sigma @ (\text{map } \text{snd } \Psi)) \preceq$   
 $(\text{map } (\text{uncurry } (\rightarrow)) \ ?\Theta @ \Gamma @ (\text{map } \text{snd } ?\Theta))$   
**proof** –  
**have**  $\text{map } (\text{uncurry } (\rightarrow)) \ \Psi \preceq \text{map } (\text{uncurry } (\rightarrow)) \ ?\Theta$   
**proof** –  
**let**  $? \Phi = \text{map } (\lambda (\psi, \sigma, \gamma). (\psi \rightarrow \gamma, \psi \rightarrow \sigma)) \ \Omega$   
**have**  $\text{map } \text{snd } ? \Phi = \text{map } (\text{uncurry } (\rightarrow)) \ \Psi$   
**using**  $\Omega(1)$  **by** *fastforce*  
**moreover have**  $\text{map } \text{fst } ? \Phi = \text{map } (\text{uncurry } (\rightarrow)) \ ?\Theta$   
**by** *fastforce*  
**hence**  $\text{mset } (\text{map } \text{fst } ? \Phi) \subseteq \# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \ ?\Theta)$   
**by** *(metis subset-mset.dual-order.refl)*  
**moreover**  
**have**  $\text{mset } (\text{map } (\lambda(\psi, \sigma, -). (\psi, \sigma)) \ \Omega) \subseteq \# \text{mset } \Psi$   
**using**  $\Omega(1)$  **by** *simp*  
**hence**  $\forall (\varphi, \chi) \in \text{set } ? \Phi. \vdash \varphi \rightarrow \chi$  **using**  $\Omega(2)$   
**proof** (*induct*  $\Omega$ )  
**case** *Nil*  
**then show** *?case by simp*  
**next**  
**case** (*Cons*  $\omega \ \Omega$ )  
**let**  $? \Phi = \text{map } (\lambda (\psi, \sigma, \gamma). (\psi \rightarrow \gamma, \psi \rightarrow \sigma)) (\omega \# \Omega)$   
**let**  $? \Phi' = \text{map } (\lambda (\psi, \sigma, \gamma). (\psi \rightarrow \gamma, \psi \rightarrow \sigma)) \ \Omega$   
**have**  $\text{mset } (\text{map } (\lambda(\psi, \sigma, -). (\psi, \sigma)) \ \Omega) \subseteq \# \text{mset } \Psi$   
 $\text{mset } (\text{map } (\lambda(-, \sigma, \gamma). (\gamma, \sigma)) \ \Omega) \subseteq \# \text{mset } \Delta$   
**using** *Cons.prem1 Cons.prem2 subset-mset.dual-order.trans* **by**  
*fastforce+*

**with** *Cons* **have**  $\forall (\varphi, \chi) \in \text{set } ?\Phi'. \vdash \varphi \rightarrow \chi$  **by** *fastforce*  
**moreover**  
**let**  $?\psi = (\lambda (\psi, -, -). \psi) \omega$   
**let**  $?\sigma = (\lambda (-, \sigma, -). \sigma) \omega$   
**let**  $? \gamma = (\lambda (-, -, \gamma). \gamma) \omega$   
**have**  $(\lambda(-, \sigma, \gamma). (\gamma, \sigma)) = (\lambda \omega. ((\lambda (-, -, \gamma). \gamma) \omega, (\lambda (-, \sigma, -). \sigma) \omega))$  **by**  
*auto*  
**hence**  $(\lambda(-, \sigma, \gamma). (\gamma, \sigma)) \omega = (? \gamma, ? \sigma)$  **by** *metis*  
**hence**  $\vdash ? \gamma \rightarrow ? \sigma$   
**using** *Cons.premis(2) mset-subset-eqD Δ(3)*  
**by** *fastforce*  
**hence**  $\vdash (? \psi \rightarrow ? \gamma) \rightarrow (? \psi \rightarrow ? \sigma)$   
**using** *modus-ponens hypothetical-syllogism*  
**by** *blast*  
**moreover have**  
 $(\lambda(\psi, \sigma, \gamma). (\psi \rightarrow \gamma, \psi \rightarrow \sigma)) =$   
 $(\lambda \omega. (((\lambda (\psi, -, -). \psi) \omega) \rightarrow ((\lambda (-, -, \gamma). \gamma) \omega),$   
 $((\lambda (\psi, -, -). \psi) \omega) \rightarrow ((\lambda (-, \sigma, -). \sigma) \omega)))$   
**by** *auto*  
**hence**  $(\lambda(\psi, \sigma, \gamma). (\psi \rightarrow \gamma, \psi \rightarrow \sigma)) \omega = ((? \psi \rightarrow ? \gamma), (? \psi \rightarrow ? \sigma))$  **by** *metis*  
**ultimately show** *?case* **by** *simp*  
**qed**  
**ultimately show** *?thesis*  
**unfolding** *stronger-theory-relation-def*  
**by** *blast*  
**qed**  
**moreover**  
**have**  $(\Sigma \ominus (\text{map snd } \Psi)) \preceq (\Gamma \ominus (\text{map snd } ?\Theta))$   
**proof** –  
**let**  $? \Delta = \Delta \ominus (\text{map } (\lambda (-, \sigma, \gamma). (\gamma, \sigma)) \Omega)$   
**have**  $\text{mset } (\text{map fst } ? \Delta) \subseteq \# \text{mset } (\Gamma \ominus (\text{map snd } ? \Theta))$   
**using**  $\Delta(2)$   
**by**  $(\text{metis } \Omega(2))$   
 $\langle \text{map snd } (\text{map } (\lambda(\psi, -, \gamma). (\psi, \gamma)) \Omega) =$   
 $\text{map fst } (\text{map } (\lambda(-, \sigma, \gamma). (\gamma, \sigma)) \Omega) \rangle$   
*list-subtract-monotonic*  
*map-list-subtract-mset-equivalence*  
**moreover**  
**from**  $\Omega(2)$  **have**  $\text{mset } ? \Delta \subseteq \# \text{mset } \Delta$  **by** *simp*  
**hence**  $\forall (\gamma, \sigma) \in \text{set } ? \Delta. \vdash \gamma \rightarrow \sigma$   
**using**  $\Delta(3)$   
**by**  $(\text{metis } \text{mset-subset-eqD set-mset-mset})$   
**moreover**  
**have**  $\text{map snd } (\text{map } (\lambda(-, \sigma, \gamma). (\gamma, \sigma)) \Omega) = \text{map snd } \Psi$   
**using**  $\Omega(1)$   
**by**  $(\text{induct } \Omega, \text{simp}, \text{fastforce})$   
**hence**  $\text{mset } (\text{map snd } ? \Delta) = \text{mset } (\Sigma \ominus (\text{map snd } \Psi))$   
**by**  $(\text{metis } \Delta(1) \Omega(2) \text{map-list-subtract-mset-equivalence})$   
**ultimately show** *?thesis*



by (*metis stronger-theory-relation-alt-def*)  
 qed  
 ultimately show *?thesis* using *stronger-theory-combine* by *blast*  
 qed  
 hence  $\text{map } (\text{uncurry } (\rightarrow)) \text{ ?}\Theta @ \Gamma \ominus (\text{map } \text{snd } \text{?}\Theta) \$\vdash \Phi$   
 using  $\Psi(3)$  *Cons* by *blast*  
 ultimately show *?case*  
 by (*metis measure-deduction.simps(2)*)  
 qed

**lemma** (in *classical-logic*) *merge-witness-measure-deduction-intro*:  
 assumes  $\text{mset } (\text{map } \text{snd } \Delta) \subseteq\# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus (\text{map } \text{snd } \Psi))$   
 and  $\text{map } (\text{uncurry } (\rightarrow)) \Delta @ (\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map } \text{snd } \Psi) \ominus \text{map } \text{snd } \Delta \$\vdash \Phi$   
 (is  $\text{?}\Gamma_0 \$\vdash \Phi$ )  
 shows  $\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{J} \Psi \Delta) @ \Gamma \ominus \text{map } \text{snd } (\mathfrak{J} \Psi \Delta) \$\vdash \Phi$   
 (is  $\text{?}\Gamma \$\vdash \Phi$ )

**proof** –  
 let  $\text{?}\Sigma = \mathfrak{B} \Psi \Delta$   
 let  $\text{?}A = \text{map } (\text{uncurry } (\rightarrow)) \Delta$   
 let  $\text{?}B = \text{map } (\text{uncurry } (\rightarrow)) \Psi$   
 let  $\text{?}C = \text{map } \text{snd } \text{?}\Sigma$   
 let  $\text{?}D = \Gamma \ominus (\text{map } \text{snd } \Psi)$   
 let  $\text{?}E = \text{map } \text{snd } (\Delta \ominus \text{?}\Sigma)$   
 have  $\Sigma: \text{mset } \text{?}\Sigma \subseteq\# \text{mset } \Delta$   
 $\text{mset } \text{?}C \subseteq\# \text{mset } \text{?}B$   
 $\text{mset } \text{?}E \subseteq\# \text{mset } \text{?}D$   
 using *assms(1)*  
*second-component-msub*  
*second-component-snd-projection-msub*  
*second-component-diff-msub*  
 by *simp+*  
**moreover**  
**from** *calculation* **have**  
 $\text{image-mset } \text{snd } (\text{mset } \Delta - \text{mset } (\mathfrak{B} \Psi \Delta))$   
 $\subseteq\# \text{mset } \Gamma - \text{image-mset } \text{snd } (\text{mset } \Psi)$   
 by *simp*  
**hence**  $\text{mset } \Gamma - \text{image-mset } \text{snd } (\text{mset } \Psi)$   
 $- \text{image-mset } \text{snd } (\text{mset } \Delta - \text{mset } (\mathfrak{B} \Psi \Delta))$   
 $+ \text{image-mset } \text{snd } (\text{mset } \Delta - \text{mset } (\mathfrak{B} \Psi \Delta))$   
 $= \text{mset } \Gamma - \text{image-mset } \text{snd } (\text{mset } \Psi)$   
 using *subset-mset.diff-add* by *blast*  
**then have**  $\text{image-mset } \text{snd } (\text{mset } \Delta - \text{mset } (\mathfrak{B} \Psi \Delta))$   
 $+ (\{\#x \rightarrow y. (x, y) \in\# \text{mset } \Psi\#})$   
 $+ (\text{mset } \Gamma - (\text{image-mset } \text{snd } (\text{mset } \Psi)$   
 $+ \text{image-mset } \text{snd } (\text{mset } \Delta - \text{mset } (\mathfrak{B} \Psi \Delta))))$   
 $= \{\#x \rightarrow y. (x, y) \in\# \text{mset } \Psi\# + (\text{mset } \Gamma - \text{image-mset } \text{snd } (\text{mset } \Psi))$   
 by (*simp add: union-commute*)

**with calculation have**  $mset \ ?\Gamma_0 = mset \ (?A @ (?B \ominus ?C) @ (?D \ominus ?E))$   
**by** (*simp, metis (no-types) add-diff-cancel-left image-mset-union subset-mset.diff-add*)  
**moreover have**  $(?A @ (?B \ominus ?C)) \preceq map \ (uncurry \ (\rightarrow)) \ (\exists \ \Psi \ \Delta)$   
**using** *second-component-merge-witness-stronger-theory* **by** *simp*  
**moreover have**  $mset \ (?D \ominus ?E) = mset \ (\Gamma \ominus map \ snd \ (\exists \ \Psi \ \Delta))$   
**using** *second-component-merge-witness-snd-projection*  
**by** *simp*  
**with calculation have**  $(?A @ (?B \ominus ?C) @ (?D \ominus ?E)) \preceq \ ?\Gamma$   
**by** (*metis*  
*(no-types, lifting)*  
*stronger-theory-combine*  
*append.assoc*  
*list-subtract-mset-homomorphism*  
*msub-stronger-theory-intro*  
*map-list-subtract-mset-containment*  
*map-list-subtract-mset-equivalence*  
*mset-subset-eq-add-right*  
*subset-mset.add-diff-inverse*  
*subset-mset.diff-add-assoc2*)  
**ultimately have**  $\ ?\Gamma_0 \preceq \ ?\Gamma$   
**unfolding** *stronger-theory-relation-alt-def*  
**by** *simp*  
**thus** *?thesis*  
**using** *assms(2) measure-stronger-theory-left-monotonic*  
**by** *blast*  
**qed**

**lemma (in classical-logic) measure-formula-right-split:**

$$\Gamma \ \$\vdash \ (\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Phi) = \Gamma \ \$\vdash \ (\varphi \# \Phi)$$

**proof** (*rule iffI*)

**assume**  $\Gamma \ \$\vdash \ (\varphi \# \Phi)$

**from this obtain**  $\Psi$  **where**  $\Psi$ :

$$mset \ (map \ snd \ \Psi) \subseteq \# \ mset \ \Gamma$$

$$map \ (uncurry \ (\sqcup)) \ \Psi \ \vdash \ \varphi$$

$$(map \ (uncurry \ (\rightarrow)) \ \Psi \ @ \ \Gamma \ \ominus \ (map \ snd \ \Psi)) \ \$\vdash \ \Phi$$

**by** *auto*

**let**  $\ ?\Psi_1 = zip \ (map \ (\lambda \ (\chi, \gamma). \ \psi \sqcup \chi) \ \Psi) \ (map \ snd \ \Psi)$

**let**  $\ ?\Gamma_1 = map \ (uncurry \ (\rightarrow)) \ \ ?\Psi_1 \ @ \ \Gamma \ \ominus \ (map \ snd \ \ ?\Psi_1)$

**let**  $\ ?\Psi_2 = zip \ (map \ (\lambda \ (\chi, \gamma). \ \psi \rightarrow \chi) \ \Psi) \ (map \ (uncurry \ (\rightarrow)) \ \ ?\Psi_1)$

**let**  $\ ?\Gamma_2 = map \ (uncurry \ (\rightarrow)) \ \ ?\Psi_2 \ @ \ \ ?\Gamma_1 \ \ominus \ (map \ snd \ \ ?\Psi_2)$

**have**  $map \ (uncurry \ (\rightarrow)) \ \Psi \preceq map \ (uncurry \ (\rightarrow)) \ \ ?\Psi_2$

**proof** (*induct*  $\Psi$ )

**case** *Nil*

**then show** *?case* **by** *simp*

**next**

**case** (*Cons*  $\delta \ \Psi$ )

**let**  $\ ?\chi = fst \ \delta$

**let**  $\ ?\gamma = snd \ \delta$

**let**  $\ ?\Psi_1 = zip \ (map \ (\lambda \ (\chi, \gamma). \ \psi \sqcup \chi) \ \Psi) \ (map \ snd \ \Psi)$

```

let ?Ψ2 = zip (map (λ (χ,γ). ψ → χ) Ψ) (map (uncurry (→)) ?Ψ1)
let ?T1 = λ Ψ. map (uncurry (→)) (zip (map (λ (χ,γ). ψ ⊔ χ) Ψ) (map snd
Ψ))
let ?T2 = λ Ψ. map (uncurry (→)) (zip (map (λ (χ,γ). ψ → χ) Ψ) (?T1 Ψ))
{
  fix δ :: 'a × 'a
  have (λ (χ,γ). ψ ⊔ χ) = (λ δ. ψ ⊔ (fst δ))
    (λ (χ,γ). ψ → χ) = (λ δ. ψ → (fst δ))
    by fastforce+
  note functional-identities = this
  have (λ (χ,γ). ψ ⊔ χ) δ = ψ ⊔ (fst δ)
    (λ (χ,γ). ψ → χ) δ = ψ → (fst δ)
    by (simp add: functional-identities)+
}
hence ?T2 (δ # Ψ) = ((ψ → ?χ) → (ψ ⊔ ?χ) → ?γ) # (map (uncurry (→))
?Ψ2)
  by simp
moreover have map (uncurry (→)) (δ # Ψ) = (?χ → ?γ) # map (uncurry
(→)) Ψ
  by (simp add: case-prod-beta)
moreover
{
  fix χ ψ γ
  have ⊢ ((ψ → χ) → (ψ ⊔ χ) → γ) ↔ (χ → γ)
  proof -
    have ∀ M. M ⊨prop ((⟨ψ⟩ → ⟨χ⟩) → (⟨ψ⟩ ⊔ ⟨χ⟩) → ⟨γ⟩) ↔ (⟨χ⟩ → ⟨γ⟩)
    by fastforce
    hence ⊢ (⟨⟨ψ⟩ → ⟨χ⟩⟩ → (⟨ψ⟩ ⊔ ⟨χ⟩) → ⟨γ⟩) ↔ (⟨χ⟩ → ⟨γ⟩) ⊔
    using propositional-semantic by blast
    thus ?thesis by simp
  qed
}
hence identity: ⊢ ((ψ → ?χ) → (ψ ⊔ ?χ) → ?γ) → (?χ → ?γ)
  using biconditional-def by auto
assume map (uncurry (→)) Ψ ≲ map (uncurry (→)) ?Ψ2
with identity have ((?χ → ?γ) # map (uncurry (→)) Ψ) ≲
  (((ψ → ?χ) → (ψ ⊔ ?χ) → ?γ) # (map (uncurry (→)) ?Ψ2))
  using stronger-theory-left-right-cons by blast
ultimately show ?case by simp
qed
hence (map (uncurry (→)) Ψ @ Γ ⊖ (map snd Ψ)) ≲
  ((map (uncurry (→)) ?Ψ2) @ Γ ⊖ (map snd Ψ))
  using stronger-theory-combine stronger-theory-reflexive by blast
moreover have mset ?T2 = mset ((map (uncurry (→)) ?Ψ2) @ Γ ⊖ (map snd
?Ψ1))
  by simp
ultimately have (map (uncurry (→)) Ψ @ Γ ⊖ (map snd Ψ)) ≲ ?T2
  by (simp add: stronger-theory-relation-def)
hence ?T2 $⊢ Φ

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    using  $\Psi(\beta)$  measure-stronger-theory-left-monotonic by blast
  moreover
  have (map (uncurry ( $\sqcup$ ))  $?\Psi_2$ )  $\vdash \psi \rightarrow \varphi$ 
  proof -
    let  $?\Gamma = \text{map } (\lambda (\chi, \gamma). (\psi \rightarrow \chi) \sqcup (\psi \sqcup \chi) \rightarrow \gamma) \Psi$ 
    let  $?\Sigma = \text{map } (\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) \Psi$ 
    have map (uncurry ( $\sqcup$ ))  $?\Psi_2 = ?\Gamma$ 
    proof (induct  $\Psi$ )
      case Nil
      then show ?case by simp
    next
      case (Cons  $\chi \Psi$ )
      have (lambda  $\varphi$ . (case  $\varphi$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \rightarrow \chi$ )  $\sqcup$  (case  $\varphi$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \sqcup \chi$ )  $\rightarrow$ 
      snd  $\varphi$ ) =
        (lambda  $\varphi$ . (case  $\varphi$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \rightarrow \chi \sqcup (\psi \sqcup \chi) \rightarrow \gamma$ ))
      by fastforce
      hence (case  $\chi$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \rightarrow \chi$ )  $\sqcup$  (case  $\chi$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \sqcup \chi$ )  $\rightarrow$  snd  $\chi$ 
      =
        (case  $\chi$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \rightarrow \chi \sqcup (\psi \sqcup \chi) \rightarrow \gamma$ )
      by metis
      with Cons show ?case by simp
    qed
  moreover have  $?\Sigma \preceq ?\Gamma$ 
  proof (induct  $\Psi$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\delta \Psi$ )
    let  $?\alpha = (\lambda (\chi, \gamma). (\psi \rightarrow \chi) \sqcup (\psi \sqcup \chi) \rightarrow \gamma) \delta$ 
    let  $?\beta = (\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) \delta$ 
    let  $?\chi = \text{fst } \delta$ 
    let  $?\gamma = \text{snd } \delta$ 
    have (lambda  $\delta$ . (case  $\delta$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \rightarrow \chi \sqcup (\psi \sqcup \chi) \rightarrow \gamma$ )) =
      (lambda  $\delta$ .  $\psi \rightarrow \text{fst } \delta \sqcup (\psi \sqcup \text{fst } \delta) \rightarrow \text{snd } \delta$ )
      (lambda  $\delta$ . (case  $\delta$  of ( $\chi, \gamma$ )  $\Rightarrow \psi \rightarrow (\chi \sqcup \gamma)$ )) = (lambda  $\delta$ .  $\psi \rightarrow (\text{fst } \delta \sqcup \text{snd } \delta)$ )
    by fastforce+
    hence  $?\alpha = (\psi \rightarrow ?\chi) \sqcup (\psi \sqcup ?\chi) \rightarrow ?\gamma$ 
       $?\beta = \psi \rightarrow (?\chi \sqcup ?\gamma)$ 
    by metis+
  moreover
  {
    fix  $\psi \chi \gamma$ 
    have  $\vdash ((\psi \rightarrow \chi) \sqcup (\psi \sqcup \chi) \rightarrow \gamma) \rightarrow (\psi \rightarrow (\chi \sqcup \gamma))$ 
    proof -
      have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} (((\psi) \rightarrow (\chi)) \sqcup ((\psi) \sqcup (\chi)) \rightarrow (\gamma)) \rightarrow ((\psi) \rightarrow ((\chi) \sqcup (\gamma)))$ 
      by fastforce
      hence  $\vdash (\Downarrow ((\psi) \rightarrow (\chi)) \sqcup ((\psi) \sqcup (\chi)) \rightarrow (\gamma)) \rightarrow ((\psi) \rightarrow ((\chi) \sqcup (\gamma))) \Downarrow$ 
      using propositional-semantics by blast
    }
  }

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    thus ?thesis by simp
  qed
}
ultimately have  $\vdash ?\alpha \rightarrow ?\beta$  by simp
thus ?case
  using Cons
    stronger-theory-left-right-cons
  by simp
qed
moreover have  $\forall \varphi. (\text{map } (\text{uncurry } (\sqcup)) \Psi) \vdash \varphi \longrightarrow ?\Sigma \vdash \psi \rightarrow \varphi$ 
proof (induct  $\Psi$ )
  case Nil
  then show ?case
    using axiom-k modus-ponens
    by fastforce
next
  case (Cons  $\delta \Psi$ )
  let ? $\delta'$  =  $(\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) \delta$ 
  let ? $\Sigma$  =  $\text{map } (\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) \Psi$ 
  let ? $\Sigma'$  =  $\text{map } (\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) (\delta \# \Psi)$ 
  {
    fix  $\varphi$ 
    assume  $\text{map } (\text{uncurry } (\sqcup)) (\delta \# \Psi) \vdash \varphi$ 
    hence  $\text{map } (\text{uncurry } (\sqcup)) \Psi \vdash (\text{uncurry } (\sqcup)) \delta \rightarrow \varphi$ 
      using list-deduction-theorem
      by simp
    hence  $?\Sigma \vdash \psi \rightarrow (\text{uncurry } (\sqcup)) \delta \rightarrow \varphi$ 
      using Cons
      by blast
    moreover
    {
      fix  $\alpha \beta \gamma$ 
      have  $\vdash (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow ((\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \gamma)$ 
        using axiom-s by auto
    }
    ultimately have  $?\Sigma \vdash (\psi \rightarrow (\text{uncurry } (\sqcup)) \delta) \rightarrow \psi \rightarrow \varphi$ 
      using list-deduction-weaken [where  $?T=?\Sigma$ ]
        list-deduction-modus-ponens [where  $?T=?\Sigma$ ]
      by metis
    moreover
    have  $(\lambda \delta. \psi \rightarrow (\text{uncurry } (\sqcup)) \delta) = (\lambda \delta. (\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) \delta)$ 
      by fastforce
    ultimately have  $?\Sigma \vdash (\lambda (\chi, \gamma). (\psi \rightarrow (\chi \sqcup \gamma))) \delta \rightarrow \psi \rightarrow \varphi$ 
      by metis
    hence  $?\Sigma' \vdash \psi \rightarrow \varphi$ 
      using list-deduction-theorem
      by simp
  }
}
then show ?case by simp

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qed
with  $\Psi(\varrho)$  have  $? \Sigma \vdash \psi \rightarrow \varphi$ 
  by blast
ultimately show ?thesis
  using stronger-theory-deduction-monotonic by auto
qed
moreover have  $mset (map\ snd\ ? \Psi_2) \subseteq \# mset\ ? \Gamma_1$  by simp
ultimately have  $? \Gamma_1 \ \$ \vdash (\psi \rightarrow \varphi \ \# \ \Phi)$  using measure-deduction.simps(2) by
blast
moreover have  $\vdash (map\ (uncurry\ (\sqcup))\ \Psi \ : \rightarrow \varphi) \rightarrow (map\ (uncurry\ (\sqcup))\ ? \Psi_1) \ : \rightarrow$ 
 $(\psi \sqcup \varphi)$ 
proof (induct  $\Psi$ )
  case Nil
  then show ?case
    unfolding disjunction-def
    using axiom-k modus-ponens
    by fastforce
next
case (Cons  $\nu\ \Psi$ )
let  $? \Delta = map\ (uncurry\ (\sqcup))\ \Psi$ 
let  $? \Delta' = map\ (uncurry\ (\sqcup))\ (\nu \ \# \ \Psi)$ 
let  $? \Sigma = map\ (uncurry\ (\sqcup))\ (zip\ (map\ (\lambda\ (\chi, \gamma). \psi \sqcup \chi)\ \Psi)\ (map\ snd\ \Psi))$ 
let  $? \Sigma' = map\ (uncurry\ (\sqcup))\ (zip\ (map\ (\lambda\ (\chi, \gamma). \psi \sqcup \chi)\ (\nu \ \# \ \Psi))\ (map\ snd\$ 
 $(\nu \ \# \ \Psi)))$ 
have  $\vdash (? \Delta' \ : \rightarrow \varphi) \rightarrow (uncurry\ (\sqcup))\ \nu \rightarrow ? \Delta \ : \rightarrow \varphi$ 
  by (simp, metis axiom-k axiom-s modus-ponens)
with Cons have  $\vdash (? \Delta' \ : \rightarrow \varphi) \rightarrow (uncurry\ (\sqcup))\ \nu \rightarrow ? \Sigma \ : \rightarrow (\psi \sqcup \varphi)$ 
  using hypothetical-syllogism modus-ponens
  by blast
hence  $(? \Delta' \ : \rightarrow \varphi) \ \# \ ((uncurry\ (\sqcup))\ \nu) \ \# \ ? \Sigma \ \vdash \psi \sqcup \varphi$ 
  by (simp add: list-deduction-def)
moreover have  $set\ ((? \Delta' \ : \rightarrow \varphi) \ \# \ ((uncurry\ (\sqcup))\ \nu) \ \# \ ? \Sigma) =$ 
 $set\ (((uncurry\ (\sqcup))\ \nu) \ \# \ (? \Delta' \ : \rightarrow \varphi) \ \# \ ? \Sigma)$ 
  by fastforce
ultimately have  $((uncurry\ (\sqcup))\ \nu) \ \# \ (? \Delta' \ : \rightarrow \varphi) \ \# \ ? \Sigma \ \vdash \psi \sqcup \varphi$ 
  using list-deduction-monotonic by blast
hence  $(? \Delta' \ : \rightarrow \varphi) \ \# \ ? \Sigma \ \vdash ((uncurry\ (\sqcup))\ \nu) \rightarrow (\psi \sqcup \varphi)$ 
  using list-deduction-theorem
  by simp
moreover
let  $? \chi = fst\ \nu$ 
let  $? \gamma = snd\ \nu$ 
have  $(\lambda\ \nu . (uncurry\ (\sqcup))\ \nu) = (\lambda\ \nu . fst\ \nu \sqcup snd\ \nu)$ 
  by fastforce
hence  $(uncurry\ (\sqcup))\ \nu = ? \chi \sqcup ? \gamma$  by simp
ultimately have  $(? \Delta' \ : \rightarrow \varphi) \ \# \ ? \Sigma \ \vdash (? \chi \sqcup ? \gamma) \rightarrow (\psi \sqcup \varphi)$  by simp
moreover
{
  fix  $\alpha\ \beta\ \delta\ \gamma$ 

```

```

    have  $\vdash ((\beta \sqcup \alpha) \rightarrow (\gamma \sqcup \delta)) \rightarrow ((\gamma \sqcup \beta) \sqcup \alpha) \rightarrow (\gamma \sqcup \delta)$ 
    proof -
      have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ((\langle \beta \rangle \sqcup \langle \alpha \rangle) \rightarrow (\langle \gamma \rangle \sqcup \langle \delta \rangle)) \rightarrow (((\langle \gamma \rangle \sqcup \langle \beta \rangle) \sqcup \langle \alpha \rangle) \rightarrow (\langle \gamma \rangle \sqcup \langle \delta \rangle))$ 
    →  $(\langle \gamma \rangle \sqcup \langle \delta \rangle)$ 
      by fastforce
      hence  $\vdash (\langle (\langle \beta \rangle \sqcup \langle \alpha \rangle) \rightarrow (\langle \gamma \rangle \sqcup \langle \delta \rangle) \rangle \rightarrow ((\langle \gamma \rangle \sqcup \langle \beta \rangle) \sqcup \langle \alpha \rangle) \rightarrow (\langle \gamma \rangle \sqcup \langle \delta \rangle))$ 
     $\Downarrow$ 
      using propositional-semantic by blast
      thus ?thesis by simp
    qed
  }
  hence  $(? \Delta' : \rightarrow \varphi) \# ? \Sigma \vdash ((? \chi \sqcup ? \gamma) \rightarrow (\psi \sqcup \varphi)) \rightarrow ((\psi \sqcup ? \chi) \sqcup ? \gamma) \rightarrow (\psi \sqcup \varphi)$ 
  using list-deduction-weaken by blast
  ultimately have  $(? \Delta' : \rightarrow \varphi) \# ? \Sigma \vdash ((\psi \sqcup ? \chi) \sqcup ? \gamma) \rightarrow (\psi \sqcup \varphi)$ 
  using list-deduction-modus-ponens by blast
  hence  $((\psi \sqcup ? \chi) \sqcup ? \gamma) \# (? \Delta' : \rightarrow \varphi) \# ? \Sigma \vdash \psi \sqcup \varphi$ 
  using list-deduction-theorem
  by simp
  moreover have  $set (((\psi \sqcup ? \chi) \sqcup ? \gamma) \# (? \Delta' : \rightarrow \varphi) \# ? \Sigma) = set ((? \Delta' : \rightarrow \varphi) \# ((\psi \sqcup ? \chi) \sqcup ? \gamma) \# ? \Sigma)$ 
  by fastforce
  moreover have
    map (uncurry  $\sqcup$ )  $(\nu \# \Psi) : \rightarrow \varphi$ 
     $\# (\psi \sqcup fst \nu) \sqcup snd \nu$ 
     $\# map (uncurry (\sqcup)) (zip (map (\lambda(-), a). \psi \sqcup a) \Psi) (map snd \Psi)) \vdash (\psi \sqcup fst \nu) \sqcup snd \nu$ 
  by (meson list.set-intros(1) list-deduction-monotonic list-deduction-reflection set-subset-Cons)
  ultimately have  $(? \Delta' : \rightarrow \varphi) \# ((\psi \sqcup ? \chi) \sqcup ? \gamma) \# ? \Sigma \vdash \psi \sqcup \varphi$ 
  using list-deduction-modus-ponens list-deduction-monotonic by blast
  moreover
  have  $(\lambda \nu. \psi \sqcup fst \nu) = (\lambda (\chi, \gamma). \psi \sqcup \chi)$ 
  by fastforce
  hence  $\psi \sqcup fst \nu = (\lambda (\chi, \gamma). \psi \sqcup \chi) \nu$ 
  by metis
  hence  $((\psi \sqcup ? \chi) \sqcup ? \gamma) \# ? \Sigma = ? \Sigma'$ 
  by simp
  ultimately have  $(? \Delta' : \rightarrow \varphi) \# ? \Sigma' \vdash \psi \sqcup \varphi$  by simp
  then show ?case by (simp add: list-deduction-def)
  qed
  with  $\Psi(2)$  have  $map (uncurry (\sqcup)) ? \Psi_1 \vdash (\psi \sqcup \varphi)$ 
  unfolding list-deduction-def
  using modus-ponens
  by blast
  moreover have  $mset (map snd ? \Psi_1) \subseteq \# mset \Gamma$  using  $\Psi(1)$  by simp
  ultimately show  $\Gamma \S \vdash (\psi \sqcup \varphi) \# \psi \rightarrow \varphi \# \Phi$ 

```

```

    using measure-deduction.simps(2) by blast
next
assume  $\Gamma \ \$\vdash (\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Phi)$ 
from this obtain  $\Psi$  where  $\Psi$ :
  mset (map snd  $\Psi$ )  $\subseteq\#$  mset  $\Gamma$ 
  map (uncurry ( $\sqcup$ ))  $\Psi \vdash \psi \sqcup \varphi$ 
  map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus$  (map snd  $\Psi$ )  $\ \$\vdash (\psi \rightarrow \varphi \# \Phi)$ 
  using measure-deduction.simps(2) by blast
let  $\ ?\Gamma' = \text{map (uncurry } (\rightarrow)) \ \Psi @ \Gamma \ominus$  (map snd  $\Psi$ )
from  $\Psi$  obtain  $\Delta$  where  $\Delta$ :
  mset (map snd  $\Delta$ )  $\subseteq\#$  mset  $\ ?\Gamma'$ 
  map (uncurry ( $\sqcup$ ))  $\Delta \vdash \psi \rightarrow \varphi$ 
  (map (uncurry ( $\rightarrow$ ))  $\Delta @ \ ?\Gamma' \ominus$  (map snd  $\Delta$ ))  $\ \$\vdash \Phi$ 
  using measure-deduction.simps(2) by blast
let  $\ ?\Omega = \ ?\ ?\Psi \ \Delta$ 
have mset (map snd  $\ ?\Omega$ )  $\subseteq\#$  mset  $\Gamma$ 
  using  $\Delta(1) \ \Psi(1)$  merge-witness-msub-intro
  by blast
moreover have map (uncurry ( $\sqcup$ ))  $\ ?\Omega \vdash \varphi$ 
proof –
  have map (uncurry ( $\sqcup$ ))  $\ ?\Omega \vdash \psi \sqcup \varphi$ 
  map (uncurry ( $\sqcup$ ))  $\ ?\Omega \vdash \psi \rightarrow \varphi$ 
  using  $\Psi(2) \ \Delta(2)$ 
  stronger-theory-deduction-monotonic
  right-merge-witness-stronger-theory
  left-merge-witness-stronger-theory
  by blast+
moreover
have  $\vdash (\psi \sqcup \varphi) \rightarrow (\psi \rightarrow \varphi) \rightarrow \varphi$ 
  unfolding disjunction-def
  using modus-ponens excluded-middle-elimination flip-implication
  by blast
ultimately show  $\ ?thesis$ 
  using list-deduction-weaken list-deduction-modus-ponens
  by blast
qed
moreover have map (uncurry ( $\rightarrow$ ))  $\ ?\Omega @ \Gamma \ominus$  (map snd  $\ ?\Omega$ )  $\ \$\vdash \Phi$ 
  using  $\Delta(1) \ \Delta(3) \ \Psi(1)$  merge-witness-measure-deduction-intro by blast
ultimately show  $\Gamma \ \$\vdash (\varphi \# \Phi)$ 
  using measure-deduction.simps(2) by blast
qed

```

**primrec (in *implication-logic*)**

*X-witness* ::  $('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \langle \mathfrak{X} \rangle$

**where**

$\mathfrak{X} \ \Psi \ [] = []$

$| \ \mathfrak{X} \ \Psi \ (\delta \# \Delta) =$

(*case find*  $(\lambda \psi. (\text{uncurry } (\rightarrow)) \ \psi = \text{snd } \delta) \ \Psi$  of

*None*  $\Rightarrow \delta \# \mathfrak{X} \ \Psi \ \Delta$



|  $\text{Some } \psi \Rightarrow (\text{fst } \psi \rightarrow \text{fst } \delta, \text{snd } \psi) \# (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta))$

**primrec** (in *implication-logic*)

$X\text{-component} :: ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} (\langle \mathfrak{X}_\bullet \rangle)$

where

$\mathfrak{X}_\bullet \Psi [] = []$

|  $\mathfrak{X}_\bullet \Psi (\delta \# \Delta) =$

(case find ( $\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta$ )  $\Psi$  of

None  $\Rightarrow \mathfrak{X}_\bullet \Psi \Delta$

|  $\text{Some } \psi \Rightarrow (\text{fst } \psi \rightarrow \text{fst } \delta, \text{snd } \psi) \# (\mathfrak{X}_\bullet (\text{remove1 } \psi \Psi) \Delta))$

**primrec** (in *implication-logic*)

$Y\text{-witness} :: ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} (\langle \mathfrak{Y} \rangle)$

where

$\mathfrak{Y} \Psi [] = \Psi$

|  $\mathfrak{Y} \Psi (\delta \# \Delta) =$

(case find ( $\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta$ )  $\Psi$  of

None  $\Rightarrow \mathfrak{Y} \Psi \Delta$

|  $\text{Some } \psi \Rightarrow (\text{fst } \psi, (\text{fst } \psi \rightarrow \text{fst } \delta) \rightarrow \text{snd } \psi) \# (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$

**primrec** (in *implication-logic*)

$Y\text{-component} :: ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} \Rightarrow ('a \times 'a) \text{ list} (\langle \mathfrak{Y}_\bullet \rangle)$

where

$\mathfrak{Y}_\bullet \Psi [] = []$

|  $\mathfrak{Y}_\bullet \Psi (\delta \# \Delta) =$

(case find ( $\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta$ )  $\Psi$  of

None  $\Rightarrow \mathfrak{Y}_\bullet \Psi \Delta$

|  $\text{Some } \psi \Rightarrow (\text{fst } \psi, (\text{fst } \psi \rightarrow \text{fst } \delta) \rightarrow \text{snd } \psi) \# (\mathfrak{Y}_\bullet (\text{remove1 } \psi \Psi) \Delta))$

**lemma** (in *implication-logic*)  $X\text{-witness-right-empty}$  [simp]:

$\mathfrak{X} [] \Delta = \Delta$

by (induct  $\Delta$ , simp+)

**lemma** (in *implication-logic*)  $Y\text{-witness-right-empty}$  [simp]:

$\mathfrak{Y} [] \Delta = []$

by (induct  $\Delta$ , simp+)

**lemma** (in *implication-logic*)  $X\text{-witness-map-snd-decomposition}$ :

$\text{mset } (\text{map } \text{snd } (\mathfrak{X} \Psi \Delta)) = \text{mset } (\text{map } \text{snd } ((\mathfrak{A} \Psi \Delta) @ (\Delta \ominus (\mathfrak{B} \Psi \Delta))))$

**proof** –

have  $\forall \Psi. \text{mset } (\text{map } \text{snd } (\mathfrak{X} \Psi \Delta)) = \text{mset } (\text{map } \text{snd } ((\mathfrak{A} \Psi \Delta) @ (\Delta \ominus (\mathfrak{B} \Psi \Delta))))$

**proof** (induct  $\Delta$ )

case Nil

then show ?case by simp

next

case (Cons  $\delta \Delta$ )

```

{
  fix  $\Psi$ 
  have mset (map snd ( $\mathfrak{X} \Psi (\delta \# \Delta)$ ))
    = mset (map snd ( $\mathfrak{A} \Psi (\delta \# \Delta) @ (\delta \# \Delta) \ominus \mathfrak{B} \Psi (\delta \# \Delta)$ ))
  using Cons
  by (cases find ( $\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{None}$ ,
      simp,
      metis (no-types, lifting)
            add-mset-add-single
            image-mset-single
            image-mset-union
            mset-subset-eq-multiset-union-diff-commute
            second-component-msub,
      fastforce)
}
then show ?case by blast
qed
thus ?thesis by blast
qed

```

**lemma** (in *implication-logic*) *Y-witness-map-snd-decomposition*:

```

mset (map snd ( $\mathfrak{Y} \Psi \Delta$ )) = mset (map snd (( $\Psi \ominus \mathfrak{A} \Psi \Delta$ ) @ ( $\mathfrak{Y} \bullet \Psi \Delta$ )))
proof -
  have  $\forall \Psi. \text{mset (map snd } (\mathfrak{Y} \Psi \Delta)) = \text{mset (map snd } ((\Psi \ominus \mathfrak{A} \Psi \Delta) @ (\mathfrak{Y} \bullet \Psi \Delta))$ 
   $\Psi \Delta))$ 
  proof (induct  $\Delta$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\delta \Delta$ )
    {
      fix  $\Psi$ 
      have mset (map snd ( $\mathfrak{Y} \Psi (\delta \# \Delta)$ )) = mset (map snd ( $\Psi \ominus \mathfrak{A} \Psi (\delta \# \Delta)$ 
  @  $\mathfrak{Y} \bullet \Psi (\delta \# \Delta)$ ))
      using Cons
      by (cases find ( $\lambda \psi. (\text{uncurry } (\rightarrow)) \psi = \text{snd } \delta) \Psi = \text{None}$ , fastforce+)
    }
    then show ?case by blast
  qed
  thus ?thesis by blast
qed

```

**lemma** (in *implication-logic*) *X-witness-msub*:

```

assumes mset (map snd  $\Psi$ )  $\subseteq\#$  mset  $\Gamma$ 
  and mset (map snd  $\Delta$ )  $\subseteq\#$  mset (map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus (\text{map snd } \Psi)$ )
shows mset (map snd ( $\mathfrak{X} \Psi \Delta$ ))  $\subseteq\#$  mset  $\Gamma$ 
proof -
  have mset (map snd ( $\Delta \ominus \mathfrak{B} \Psi \Delta$ ))  $\subseteq\#$  mset ( $\Gamma \ominus (\text{map snd } \Psi)$ )

```

**using** *assms second-component-diff-msub* **by** *blast*  
**moreover have**  $mset (map\ snd\ (\mathfrak{A}\ \Psi\ \Delta)) \subseteq\# mset (map\ snd\ \Psi)$   
**using** *first-component-msub*  
**by** (*simp add: image-mset-subseteq-mono*)  
**moreover have**  $mset ((map\ snd\ \Psi) @ (\Gamma \ominus map\ snd\ \Psi)) = mset\ \Gamma$   
**using** *assms(1)*  
**by** *simp*  
**moreover have**  $image\ mset\ snd (mset (\mathfrak{A}\ \Psi\ \Delta)) + image\ mset\ snd (mset (\Delta \ominus \mathfrak{B}\ \Psi\ \Delta))$   
 $= mset (map\ snd (\mathfrak{X}\ \Psi\ \Delta))$   
**using** *X-witness-map-snd-decomposition* **by** *force*  
**ultimately**  
**show** *?thesis*  
**by** (*metis (no-types) mset-append mset-map subset-mset.add-mono*)  
**qed**

**lemma** (*in implication-logic*) *Y-component-msub*:

$mset (map\ snd (\mathfrak{Y}\bullet\ \Psi\ \Delta)) \subseteq\# mset (map (uncurry (\rightarrow)) (\mathfrak{X}\ \Psi\ \Delta))$   
**proof** –  
**have**  $\forall \Psi. mset (map\ snd (\mathfrak{Y}\bullet\ \Psi\ \Delta)) \subseteq\# mset (map (uncurry (\rightarrow)) (\mathfrak{X}\ \Psi\ \Delta))$   
**proof** (*induct*  $\Delta$ )  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons*  $\delta\ \Delta$ )  
{  
**fix**  $\Psi$   
**have**  $mset (map\ snd (\mathfrak{Y}\bullet\ \Psi (\delta\ \# \Delta))) \subseteq\# mset (map (uncurry (\rightarrow)) (\mathfrak{X}\ \Psi (\delta\ \# \Delta)))$   
**using** *Cons*  
**by** (*cases find* ( $\lambda \psi. (uncurry (\rightarrow)) \psi = snd\ \delta$ )  $\Psi = None$ ,  
*simp, metis add-mset-add-single*  
*mset-subset-eq-add-left*  
*subset-mset.order-trans,*  
*fastforce*)  
}  
**then show** *?case* **by** *blast*  
**qed**  
**thus** *?thesis* **by** *blast*  
**qed**

**lemma** (*in implication-logic*) *Y-witness-msub*:

**assumes**  $mset (map\ snd\ \Psi) \subseteq\# mset\ \Gamma$   
**and**  $mset (map\ snd\ \Delta) \subseteq\# mset (map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus (map\ snd\ \Psi))$   
**shows**  $mset (map\ snd (\mathfrak{Y}\ \Psi\ \Delta)) \subseteq\#$   
 $mset (map (uncurry (\rightarrow)) (\mathfrak{X}\ \Psi\ \Delta) @ \Gamma \ominus map\ snd (\mathfrak{X}\ \Psi\ \Delta))$   
**proof** –  
**have** *A*:  $image\ mset\ snd (mset\ \Psi) \subseteq\# mset\ \Gamma$  **using** *assms* **by** *simp*

**have**  $B$ : *image-mset snd* ( $mset (\mathfrak{A} \Psi \Delta)$ ) + *image-mset snd* ( $mset \Delta - mset (\mathfrak{B} \Psi \Delta)$ )  $\subseteq\#$   $mset \Gamma$   
**using**  $A$  *X-witness-map-snd-decomposition* *assms(2)* *X-witness-msub* **by** *auto*  
**have**  $mset \Gamma - image-mset\ snd (mset \Psi) = mset (\Gamma \ominus map\ snd \Psi)$   
**by** *simp*  
**then have**  $C$ :  $mset (map\ snd (\Delta \ominus \mathfrak{B} \Psi \Delta)) + image-mset\ snd (mset \Psi) \subseteq\#$   
 $mset \Gamma$   
**using**  $A$  **by** (*metis* (*full-types*) *assms(2)* *second-component-diff-msub* *subset-mset.le-diff-conv2*)  
**have** *image-mset snd* ( $mset (\Psi \ominus \mathfrak{A} \Psi \Delta)$ ) + *image-mset snd* ( $mset (\mathfrak{A} \Psi \Delta)$ )  
 $= image-mset\ snd (mset \Psi)$   
**by** (*metis* (*no-types*) *image-mset-union*  
*list-subtract-mset-homomorphism*  
*first-component-msub*  
*subset-mset.diff-add*)  
**then have** *image-mset snd* ( $mset \Psi - mset (\mathfrak{A} \Psi \Delta)$ )  
+ (*image-mset snd* ( $mset (\mathfrak{A} \Psi \Delta)$ ) + *image-mset snd* ( $mset \Delta - mset$   
 $(\mathfrak{B} \Psi \Delta)$ ))  
 $= mset (map\ snd (\Delta \ominus \mathfrak{B} \Psi \Delta)) + image-mset\ snd (mset \Psi)$   
**by** (*simp add: union-commute*)  
**then have** *image-mset snd* ( $mset \Psi - mset (\mathfrak{A} \Psi \Delta)$ )  
 $\subseteq\# mset \Gamma - (image-mset\ snd (mset (\mathfrak{A} \Psi \Delta)) + image-mset\ snd (mset$   
 $\Delta - mset (\mathfrak{B} \Psi \Delta)))$   
**by** (*metis* (*no-types*)  $B$   $C$  *subset-mset.le-diff-conv2*)  
**hence**  $mset (map\ snd (\Psi \ominus \mathfrak{A} \Psi \Delta)) \subseteq\# mset (\Gamma \ominus map\ snd (\mathfrak{X} \Psi \Delta))$   
**using** *assms* *X-witness-map-snd-decomposition*  
**by** *simp*  
**thus** *?thesis*  
**using** *Y-component-msub*  
*Y-witness-map-snd-decomposition*  
**by** (*simp add: mset-subset-eq-mono-add union-commute*)  
**qed**

**lemma** (*in classical-logic*) *X-witness-right-stronger-theory*:

*map (uncurry (□)) Δ ≲ map (uncurry (□)) (⊗ Ψ Δ)*

**proof** –

**have**  $\forall \Psi. map (uncurry (\square)) \Delta \preceq map (uncurry (\square)) (\otimes \Psi \Delta)$

**proof** (*induct Δ*)

**case** *Nil*

**then show** *?case by simp*

**next**

**case** (*Cons δ Δ*)

{

**fix**  $\Psi$

**have**  $map (uncurry (\square)) (\delta \# \Delta) \preceq map (uncurry (\square)) (\otimes \Psi (\delta \# \Delta))$

**proof** (*cases find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None*)

**case** *True*

**then show** *?thesis*

**using** *Cons*

```

    by (simp add: stronger-theory-left-right-cons
        trivial-implication)
next
case False
from this obtain  $\psi$  where
 $\psi$ : find ( $\lambda\psi$ . uncurry ( $\rightarrow$ )  $\psi = \text{snd } \delta$ )  $\Psi = \text{Some } \psi$ 
 $\psi \in \text{set } \Psi$ 
( $\text{fst } \psi \rightarrow \text{snd } \psi$ ) =  $\text{snd } \delta$ 
using find-Some-set-membership
find-Some-predicate
by fastforce
let  $?\Psi' = \text{remove1 } \psi \Psi$ 
let  $?\alpha = \text{fst } \psi$ 
let  $?\beta = \text{snd } \psi$ 
let  $?\gamma = \text{fst } \delta$ 
have map (uncurry ( $\sqcup$ ))  $\Delta \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{X} ?\Psi' \Delta)$ 
using Cons by simp
moreover
have (uncurry ( $\sqcup$ )) = ( $\lambda \delta$ .  $\text{fst } \delta \sqcup \text{snd } \delta$ ) by fastforce
hence (uncurry ( $\sqcup$ ))  $\delta = ?\gamma \sqcup (?\alpha \rightarrow ?\beta)$  using  $\psi(\beta)$  by fastforce
moreover
{
fix  $\alpha \beta \gamma$ 
have  $\vdash (\alpha \rightarrow \gamma \sqcup \beta) \rightarrow (\gamma \sqcup (\alpha \rightarrow \beta))$ 
proof -
let  $?\varphi = (\langle \alpha \rangle \rightarrow \langle \gamma \rangle \sqcup \langle \beta \rangle) \rightarrow (\langle \gamma \rangle \sqcup (\langle \alpha \rangle \rightarrow \langle \beta \rangle))$ 
have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
hence  $\vdash (\langle ?\varphi \rangle)$  using propositional-semantic by blast
thus  $?thesis$  by simp
qed
}
hence  $\vdash (?\alpha \rightarrow ?\gamma \sqcup ?\beta) \rightarrow (?\gamma \sqcup (?\alpha \rightarrow ?\beta))$  by simp
ultimately
show  $?thesis$  using  $\psi$ 
by (simp add: stronger-theory-left-right-cons)
qed
}
then show  $?case$  by simp
qed
thus  $?thesis$  by simp
qed

lemma (in classical-logic) Y-witness-left-stronger-theory:
map (uncurry ( $\sqcup$ ))  $\Psi \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{Y} \Psi \Delta)$ 
proof -
have  $\forall \Psi$ . map (uncurry ( $\sqcup$ ))  $\Psi \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{Y} \Psi \Delta)$ 
proof (induct  $\Delta$ )
case Nil
then show  $?case$  by simp

```

```

next
case (Cons δ Δ)
{
  fix Ψ
  have map (uncurry (⊔)) Ψ ≤ map (uncurry (⊔)) (ℑ Ψ (δ # Δ))
  proof (cases find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None)
    case True
    then show ?thesis using Cons by simp
  next
  case False
  from this obtain ψ where
    ψ: find (λψ. uncurry (→) ψ = snd δ) Ψ = Some ψ
    ψ ∈ set Ψ
    (uncurry (⊔)) ψ = fst ψ ⊔ snd ψ
  using find-Some-set-membership
  by fastforce
  let ?φ = fst ψ ⊔ (fst ψ → fst δ) → snd ψ
  let ?Ψ' = remove1 ψ Ψ
  have map (uncurry (⊔)) ?Ψ' ≤ map (uncurry (⊔)) (ℑ ?Ψ' Δ)
    using Cons by simp
  moreover
  {
    fix α β γ
    have ⊢ (α ⊔ (α → γ) → β) → (α ⊔ β)
    proof -
      let ?φ = ((α) ⊔ ((α) → (γ)) → (β)) → ((α) ⊔ (β))
      have ∀ℳ. ℳ ⊨prop ?φ by fastforce
      hence ⊢ (⊥ ?φ ⊤) using propositional-semantic by blast
      thus ?thesis by simp
    qed
  }
  hence ⊢ ?φ → (uncurry (⊔)) ψ using ψ(β) by auto
  ultimately
  have map (uncurry (⊔)) (ψ # ?Ψ') ≤ (?φ # map (uncurry (⊔)) (ℑ ?Ψ'
Δ))
    by (simp add: stronger-theory-left-right-cons)
  moreover
  from ψ have mset (map (uncurry (⊔)) (ψ # ?Ψ')) = mset (map (uncurry
(⊔)) Ψ)
    by (metis mset-map perm-remove)
  ultimately show ?thesis
    using stronger-theory-relation-alt-def ψ(1) by auto
  qed
}
then show ?case by blast
qed
thus ?thesis by blast
qed

```

**lemma** (in *implication-logic*) *X-witness-second-component-diff-decomposition*:  
 $mset (\mathfrak{X} \Psi \Delta) = mset (\mathfrak{X}_\bullet \Psi \Delta @ \Delta \ominus \mathfrak{B} \Psi \Delta)$   
**proof** –  
**have**  $\forall \Psi. mset (\mathfrak{X} \Psi \Delta) = mset (\mathfrak{X}_\bullet \Psi \Delta @ \Delta \ominus \mathfrak{B} \Psi \Delta)$   
**proof** (*induct*  $\Delta$ )  
  **case** *Nil*  
  **then show** *?case by simp*  
**next**  
  **case** (*Cons*  $\delta \Delta$ )  
  {  
  **fix**  $\Psi$   
  **have**  $mset (\mathfrak{X} \Psi (\delta \# \Delta)) =$   
   $mset (\mathfrak{X}_\bullet \Psi (\delta \# \Delta) @ (\delta \# \Delta) \ominus \mathfrak{B} \Psi (\delta \# \Delta))$   
  **using** *Cons*  
  **by** (*cases find* ( $\lambda \psi. (uncurry (\rightarrow)) \psi = snd \delta$ )  $\Psi = None$ ,  
  *simp*, *metis add-mset-add-single second-component-msub subset-mset.diff-add-assoc2*,  
  *fastforce*)  
  }  
  **then show** *?case by blast*  
**qed**  
**thus** *?thesis by blast*  
**qed**

**lemma** (in *implication-logic*) *Y-witness-first-component-diff-decomposition*:  
 $mset (\mathfrak{Y} \Psi \Delta) = mset (\Psi \ominus \mathfrak{A} \Psi \Delta @ \mathfrak{Y}_\bullet \Psi \Delta)$   
**proof** –  
**have**  $\forall \Psi. mset (\mathfrak{Y} \Psi \Delta) = mset (\Psi \ominus \mathfrak{A} \Psi \Delta @ \mathfrak{Y}_\bullet \Psi \Delta)$   
**proof** (*induct*  $\Delta$ )  
  **case** *Nil*  
  **then show** *?case by simp*  
**next**  
  **case** (*Cons*  $\delta \Delta$ )  
  {  
  **fix**  $\Psi$   
  **have**  $mset (\mathfrak{Y} \Psi (\delta \# \Delta)) =$   
   $mset (\Psi \ominus \mathfrak{A} \Psi (\delta \# \Delta) @ \mathfrak{Y}_\bullet \Psi (\delta \# \Delta))$   
  **using** *Cons*  
  **by** (*cases find* ( $\lambda \psi. (uncurry (\rightarrow)) \psi = snd \delta$ )  $\Psi = None$ , *simp*, *fastforce*)  
  }  
  **then show** *?case by blast*  
**qed**  
**thus** *?thesis by blast*  
**qed**

**lemma** (in *implication-logic*) *Y-witness-right-stronger-theory*:  
 $map (uncurry (\rightarrow)) \Delta \preceq map (uncurry (\rightarrow)) (\mathfrak{Y} \Psi \Delta \ominus (\Psi \ominus \mathfrak{A} \Psi \Delta) @ (\Delta \ominus \mathfrak{B} \Psi \Delta))$   
**proof** –  
  **let**  $\mathfrak{F} = \lambda \Psi \Delta. (\Psi \ominus \mathfrak{A} \Psi \Delta)$

```

let ?g = λ Ψ Δ. (Δ ⊖ ℑ Ψ Δ)
have ∀ Ψ. map (uncurry (→)) Δ ≼ map (uncurry (→)) (ℑ Ψ Δ ⊖ ?f Ψ Δ @
?g Ψ Δ)
proof (induct Δ)
  case Nil
  then show ?case by simp
next
case (Cons δ Δ)
let ?δ = (uncurry (→)) δ
{
  fix Ψ
  have map (uncurry (→)) (δ # Δ)
    ≼ map (uncurry (→)) (ℑ Ψ (δ # Δ) ⊖ ?f Ψ (δ # Δ) @ ?g Ψ (δ # Δ))
  proof (cases find (λ ψ. (uncurry (→)) ψ = snd δ) Ψ = None)
  case True
  moreover
  from Cons have
    map (uncurry (→)) (δ # Δ) ≼ map (uncurry (→)) (δ # ℑ Ψ Δ ⊖ ?f Ψ
Δ @ ?g Ψ Δ)
  by (simp add: stronger-theory-left-right-cons trivial-implication)
  moreover
  have mset (map (uncurry (→)) (δ # ℑ Ψ Δ ⊖ ?f Ψ Δ @ ?g Ψ Δ))
    = mset (map (uncurry (→)) (ℑ Ψ Δ ⊖ ?f Ψ Δ @ ((δ # Δ) ⊖ ℑ Ψ Δ)))
  by (simp,
      metis (no-types, lifting)
            add-mset-add-single
            image-mset-single
            image-mset-union
            second-component-msub
            mset-subset-eq-multiset-union-diff-commute)
  moreover have
    ∀ Ψ Φ. Ψ ≼ Φ
    = (∃ Σ. map snd Σ = Ψ
      ∧ mset (map fst Σ) ⊆# mset Φ
      ∧ (∀ ξ. ξ ∉ set Σ ∨ ⊢ (uncurry (→) ξ)))
    by (simp add: Ball-def-raw stronger-theory-relation-def)
  moreover have
    ((uncurry (→) δ) # map (uncurry (→)) Δ)
    ≼ ((uncurry (→) δ) # map (uncurry (→)) (ℑ Ψ Δ ⊖ (?f Ψ Δ))
      @ map (uncurry (→)) (?g Ψ Δ))
  using calculation by auto
  ultimately show ?thesis
  by (simp, metis union-mset-add-mset-right)
}
next
case False
from this obtain ψ where
  ψ: find (λ ψ. uncurry (→) ψ = snd δ) Ψ = Some ψ
      uncurry (→) ψ = snd δ
using find-Some-predicate

```



by *fastforce*  
 let  $?\alpha = \text{fst } \psi$   
 let  $?\beta = \text{fst } \delta$   
 let  $? \gamma = \text{snd } \psi$   
 have  $(\lambda \delta. \text{fst } \delta \rightarrow \text{snd } \delta) = \text{uncurry } (\rightarrow)$  by *fastforce*  
 hence  $? \beta \rightarrow ? \alpha \rightarrow ? \gamma = \text{uncurry } (\rightarrow) \delta$  using  $\psi(2)$  by *metis*  
 moreover  
 let  $?A = \mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta$   
 let  $?B = \mathfrak{X} (\text{remove1 } \psi \Psi) \Delta$   
 let  $?C = \mathfrak{B} (\text{remove1 } \psi \Psi) \Delta$   
 let  $?D = ?A \ominus ((\text{remove1 } \psi \Psi) \ominus ?B)$   
 have  $\text{mset } ((\text{remove1 } \psi \Psi) \ominus ?B) \subseteq \# \text{mset } ?A$   
 using *Y-witness-first-component-diff-decomposition* by *simp*  
 {  
 assume  $\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)) \subseteq \# \text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta)$   
 moreover have  $B: \forall \Phi \Psi. \exists \Delta. \Psi \subseteq \# \Phi \longrightarrow \Psi + \Delta = \Phi$   
 by *(metis subset-mset.le-iff-add)*  
 moreover obtain *f* where  
 $A: \text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta)$   
 $\quad - (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 $\quad = f (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad \quad (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 by *blast*  
 ultimately obtain *g* where  
 $B: \forall p. \text{add-mset } p (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad - (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 $\quad = \text{add-mset } p$   
 $\quad \quad (g (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad \quad (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta))))$   
 by *(metis add-diff-cancel-left' union-mset-add-mset-right)*  
 have  $g (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 $\quad = \text{add-mset } (\text{fst } \psi, (\text{fst } \psi \rightarrow \text{fst } \delta) \rightarrow \text{snd } \psi)$   
 $\quad \quad (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad - (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 $\quad - \{ \# (\text{fst } \psi, (\text{fst } \psi \rightarrow \text{fst } \delta) \rightarrow \text{snd } \psi) \# \}$   
 by *(simp add: B)*  
 then have *C*:  
 $g (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 $\quad = \text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta)$   
 $\quad \quad - (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$   
 by *simp*  
 let  $?S_1 =$   
 $\{ \# x \rightarrow y.$   
 $\quad (x, y) \in \# \text{add-mset } (\text{fst } \psi, (\text{fst } \psi \rightarrow \text{fst } \delta) \rightarrow \text{snd } \psi)$   
 $\quad \quad (\text{mset } (\mathfrak{Y} (\text{remove1 } \psi \Psi) \Delta))$   
 $\quad \quad - (\text{mset } \Psi - \text{add-mset } \psi (\text{mset } (\mathfrak{X} (\text{remove1 } \psi \Psi) \Delta)))$

```

#}
let ?S2 =
  add-mset
  (fst ψ → (fst ψ → fst δ) → snd ψ)
  {# x → y.
    (x, y) ∈# mset (ℳ) (remove1 ψ Ψ) Δ)
    - (mset Ψ
      - add-mset ψ (mset (ℳ) (remove1 ψ Ψ) Δ)))
#}
have ?S1 = ?S2
  using A C by (simp add: B)
}
hence mset (map (uncurry (→)))
  (((?α, (?α → ?β) → ?γ) # ?A) ⊖ remove1 ψ (Ψ ⊖ ?B)
   @ (remove1 δ ((δ # Δ) ⊖ ?C)))
  = mset ((?α → (?α → ?β) → ?γ) # map (uncurry (→)) (?D @ (Δ
⊖ ?C)))
  using
  add-mset-add-single
  image-mset-add-mset
  prod.simps(2)
  subset-mset.diff-add-assoc2
  ⟨mset (remove1 ψ Ψ ⊖ ℳ) (remove1 ψ Ψ) Δ⟩ ⊆# mset (ℳ) (remove1 ψ
Ψ) Δ⟩
  by fastforce
moreover
have ⊢ (?α → (?α → ?β) → ?γ) → ?β → ?α → ?γ
proof -
  let ?Γ = [(?α → (?α → ?β) → ?γ), ?β, ?α]
  have ?Γ ⊢ ?α → (?α → ?β) → ?γ
    ?Γ ⊢ ?α
  by (simp add: list-deduction-reflection)+
  hence ?Γ ⊢ (?α → ?β) → ?γ
  using list-deduction-modus-ponens by blast
  moreover have ?Γ ⊢ ?β
  by (simp add: list-deduction-reflection)
  hence ?Γ ⊢ ?α → ?β
  using axiom-k list-deduction-modus-ponens list-deduction-weaken by blast
  ultimately have ?Γ ⊢ ?γ
  using list-deduction-modus-ponens by blast
  thus ?thesis
  unfolding list-deduction-def by simp
qed
hence (?β → ?α → ?γ # map (uncurry (→)) Δ) ⋖
  (?α → (?α → ?β) → ?γ # map (uncurry (→)) (?D @ (Δ ⊖ ?C)))
  using Cons stronger-theory-left-right-cons by blast
  ultimately show ?thesis
  using ψ by (simp add: stronger-theory-relation-alt-def)
qed

```

```

    }
    then show ?case by blast
qed
thus ?thesis by blast
qed

```

**lemma** (in *implication-logic*) *xcomponent-ycomponent-connection*:

$map (uncurry (\rightarrow)) (\mathfrak{X}\bullet \Psi \Delta) = map\ snd (\mathfrak{Y}\bullet \Psi \Delta)$

**proof** –

**have**  $\forall \Psi. map (uncurry (\rightarrow)) (\mathfrak{X}\bullet \Psi \Delta) = map\ snd (\mathfrak{Y}\bullet \Psi \Delta)$

**proof** (*induct*  $\Delta$ )

**case** *Nil*

**then show** ?case by *simp*

**next**

**case** (*Cons*  $\delta \Delta$ )

    {

**fix**  $\Psi$

**have**  $map (uncurry (\rightarrow)) (\mathfrak{X}\bullet \Psi (\delta \# \Delta)) = map\ snd (\mathfrak{Y}\bullet \Psi (\delta \# \Delta))$

**using** *Cons*

**by** (*cases find*  $(\lambda \psi. (uncurry (\rightarrow)) \psi = snd \delta) \Psi = None, simp, fastforce$ )

    }

**then show** ?case by *blast*

**qed**

**thus** ?thesis by *blast*

**qed**

**lemma** (in *classical-logic*) *xwitness-ywitness-measure-deduction-intro*:

**assumes**  $mset (map\ snd \Psi) \subseteq\# mset \Gamma$

**and**  $mset (map\ snd \Delta) \subseteq\# mset (map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus (map\ snd \Psi))$

**and**  $map (uncurry (\rightarrow)) \Delta @ (map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus map\ snd \Psi) \ominus map\ snd \Delta \S\vdash \Phi$

**(is**  $? \Gamma_0 \S\vdash \Phi$ )

**shows**  $map (uncurry (\rightarrow)) (\mathfrak{Y} \Psi \Delta) @$

$(map (uncurry (\rightarrow)) (\mathfrak{X} \Psi \Delta) @ \Gamma \ominus map\ snd (\mathfrak{X} \Psi \Delta)) \ominus$

$map\ snd (\mathfrak{Y} \Psi \Delta) \S\vdash \Phi$

**(is**  $? \Gamma \S\vdash \Phi$ )

**proof** –

**let**  $?A = map (uncurry (\rightarrow)) (\mathfrak{Y} \Psi \Delta)$

**let**  $?B = map (uncurry (\rightarrow)) (\mathfrak{X} \Psi \Delta)$

**let**  $?C = \Psi \ominus \mathfrak{A} \Psi \Delta$

**let**  $?D = map (uncurry (\rightarrow)) ?C$

**let**  $?E = \Delta \ominus \mathfrak{B} \Psi \Delta$

**let**  $?F = map (uncurry (\rightarrow)) ?E$

**let**  $?G = map\ snd (\mathfrak{B} \Psi \Delta)$

**let**  $?H = map (uncurry (\rightarrow)) (\mathfrak{X}\bullet \Psi \Delta)$

**let**  $?I = \mathfrak{A} \Psi \Delta$

**let**  $?J = map\ snd (\mathfrak{X} \Psi \Delta)$

**let**  $?K = map\ snd (\mathfrak{Y} \Psi \Delta)$

**have**  $mset (map (uncurry (\rightarrow)) (\mathfrak{Y} \Psi \Delta \ominus ?C @ ?E)) = mset (?A \ominus ?D @ ?F)$   
**by** (*simp add: Y-witness-first-component-diff-decomposition*)  
**hence**  $(map (uncurry (\rightarrow)) \Delta) \preceq (?A \ominus ?D @ ?F)$   
**using** *Y-witness-right-stronger-theory*  
*stronger-theory-relation-alt-def*  
**by** (*simp, metis (no-types, lifting)*)  
**hence**  $? \Gamma_0 \preceq ((?A \ominus ?D @ ?F) @ (map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus map\ snd \Psi))$   
 $\ominus map\ snd \Delta)$   
**using** *stronger-theory-combine stronger-theory-reflexive* **by** *blast*  
**moreover**  
**have**  $\spadesuit$ :  $mset ?G \subseteq\# mset (map (uncurry (\rightarrow)) \Psi)$   
 $mset (\mathfrak{B} \Psi \Delta) \subseteq\# mset \Delta$   
 $mset (map\ snd ?E) \subseteq\# mset (\Gamma \ominus map\ snd \Psi)$   
 $mset (map (uncurry (\rightarrow)) \Psi \ominus ?G) = mset ?D$   
 $mset ?D \subseteq\# mset ?A$   
 $mset (map\ snd ?I) \subseteq\# mset (map\ snd \Psi)$   
 $mset (map\ snd ?I) \subseteq\# mset \Gamma$   
 $mset (map\ snd (?I @ ?E)) = mset ?J$   
**using** *second-component-msub*  
*second-component-diff-msub*  
*second-component-snd-projection-msub*  
*first-component-second-component-mset-connection*  
*X-witness-map-snd-decomposition*  
  
**by** (*simp,*  
*simp,*  
*metis assms(2),*  
*simp add: image-mset-Diff first-component-msub,*  
*simp add: Y-witness-first-component-diff-decomposition,*  
*simp add: image-mset-subseteq-mono first-component-msub,*  
*metis assms(1) first-component-msub map-monotonic subset-mset.dual-order.trans,*  
*simp*)  
**hence**  $mset \Delta - mset (\mathfrak{B} \Psi \Delta) + mset (\mathfrak{B} \Psi \Delta) = mset \Delta$   
**by** *simp*  
**hence**  $\heartsuit$ :  $\{\#x \rightarrow y. (x, y) \in\# mset \Psi\# \} + (mset \Gamma - image\ mset\ snd (mset$   
 $\Psi))$   

$$- image\ mset\ snd (mset \Delta)$$

$$= \{\#x \rightarrow y. (x, y) \in\# mset \Psi\# \} + (mset \Gamma - image\ mset\ snd (mset$$
  
 $\Psi))$   

$$- image\ mset\ snd (mset \Delta - mset (\mathfrak{B} \Psi \Delta))$$

$$- image\ mset\ snd (mset (\mathfrak{B} \Psi \Delta))$$

$$image\ mset\ snd (mset \Psi - mset (\mathfrak{A} \Psi \Delta)) + image\ mset\ snd (mset (\mathfrak{A}$$
  
 $\Psi \Delta))$   

$$= image\ mset\ snd (mset \Psi)$$
  
**using**  $\spadesuit$   
**by** (*metis (no-types) diff-diff-add-mset image-mset-union,*  
*metis (no-types) image-mset-union first-component-msub subset-mset.diff-add*)  
**then have**  $mset \Gamma - image\ mset\ snd (mset \Psi)$   
 $- image\ mset\ snd (mset \Delta - mset (\mathfrak{B} \Psi \Delta))$

$= \text{mset } \Gamma - (\text{image-mset snd } (\text{mset } \Psi - \text{mset } (\mathfrak{A} \Psi \Delta))$   
 $\quad + \text{image-mset snd } (\text{mset } (\mathfrak{X} \Psi \Delta)))$   
**using**  $\spadesuit$  **by** (*simp*, *metis* (*full-types*) *diff-diff-add-mset*)  
**hence**  $\text{mset } ((\text{map } (\text{uncurry } (\rightarrow)) \Psi @ \Gamma \ominus \text{map snd } \Psi) \ominus \text{map snd } \Delta)$   
 $= \text{mset } (?D @ (\Gamma \ominus ?J) \ominus \text{map snd } ?C)$   
**using**  $\heartsuit \spadesuit$  **by** (*simp*, *metis* (*no-types*) *add.commute subset-mset.add-diff-assoc*)  
**ultimately have**  $?T_0 \preceq ((?A \ominus ?D @ ?F) @ ?D @ (\Gamma \ominus ?J) \ominus \text{map snd } ?C)$   
**unfolding** *stronger-theory-relation-alt-def*  
**by** *simp*  
**moreover**  
**have**  $\text{mset } ?F = \text{mset } (?B \ominus ?H)$   
 $\text{mset } ?D \subseteq \# \text{mset } ?A$   
 $\text{mset } (\text{map snd } (\Psi \ominus ?I)) \subseteq \# \text{mset } (\Gamma \ominus ?J)$   
**by** (*simp add: X-witness-second-component-diff-decomposition*,  
*simp add: Y-witness-first-component-diff-decomposition*,  
*simp*, *metis* (*no-types*, *lifting*)  
 $\heartsuit(2) \spadesuit(8)$  *add.assoc assms(1) assms(2) image-mset-union*  
*X-witness-msub merge-witness-msub-intro*  
*second-component-merge-witness-snd-projection*  
*mset-map*  
*subset-mset.le-diff-conv2*  
*union-code*)  
**hence**  $\text{mset } ((?A \ominus ?D @ ?F) @ ?D @ (\Gamma \ominus ?J) \ominus \text{map snd } ?C)$   
 $= \text{mset } (?A @ (?B \ominus ?H @ \Gamma \ominus ?J) \ominus \text{map snd } ?C)$   
 $\text{mset } ?H \subseteq \# \text{mset } ?B$   
 $\{\#x \rightarrow y. (x, y) \in \# \text{mset } (\mathfrak{X}_\bullet \Psi \Delta)\# \} = \text{mset } (\text{map snd } (\mathfrak{Y}_\bullet \Psi \Delta))$   
**by** (*simp add: subset-mset.diff-add-assoc*,  
*simp add: X-witness-second-component-diff-decomposition*,  
*metis xcomponent-ycomponent-connection mset-map uncurry-def*)  
**hence**  $\text{mset } ((?A \ominus ?D @ ?F) @ ?D @ (\Gamma \ominus ?J) \ominus \text{map snd } ?C)$   
 $= \text{mset } (?A @ (?B @ \Gamma \ominus ?J) \ominus (?H @ \text{map snd } ?C))$   
 $\{\#x \rightarrow y. (x, y) \in \# \text{mset } (\mathfrak{X}_\bullet \Psi \Delta)\# \} + \text{image-mset snd } (\text{mset } \Psi - \text{mset } (\mathfrak{A} \Psi \Delta))$   
 $= \text{mset } (\text{map snd } (\mathfrak{Y} \Psi \Delta))$   
**using** *Y-witness-map-snd-decomposition*  
**by** (*simp add: subset-mset.diff-add-assoc*, *force*)  
**hence**  $\text{mset } ((?A \ominus ?D @ ?F) @ ?D @ (\Gamma \ominus ?J) \ominus \text{map snd } ?C)$   
 $= \text{mset } (?A @ (?B @ \Gamma \ominus ?J) \ominus ?K)$   
**by** (*simp*)  
**ultimately have**  $?T_0 \preceq (?A @ (?B @ \Gamma \ominus ?J) \ominus ?K)$   
**unfolding** *stronger-theory-relation-alt-def*  
**by** *metis*  
**thus** *?thesis*  
**using** *assms(3) measure-stronger-theory-left-monotonic*  
**by** *blast*  
**qed**

**lemma** (in *classical-logic*) *measure-cons-cons-right-permute*:  
**assumes**  $\Gamma \ \$\vdash (\varphi \# \psi \# \Phi)$

**shows**  $\Gamma \text{\$}\vdash (\psi \# \varphi \# \Phi)$   
**proof** –  
**from** *assms* **obtain**  $\Psi$  **where**  $\Psi$ :  
 $mset (map \text{snd } \Psi) \subseteq \# mset \Gamma$   
 $map (uncurry (\sqcup)) \Psi \text{\$}\vdash \varphi$   
 $map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus (map \text{snd } \Psi) \text{\$}\vdash (\psi \# \Phi)$   
**by** *fastforce*  
**let**  $?T_0 = map (uncurry (\rightarrow)) \Psi @ \Gamma \ominus (map \text{snd } \Psi)$   
**from**  $\Psi(3)$  **obtain**  $\Delta$  **where**  $\Delta$ :  
 $mset (map \text{snd } \Delta) \subseteq \# mset ?T_0$   
 $map (uncurry (\sqcup)) \Delta \text{\$}\vdash \psi$   
 $(map (uncurry (\rightarrow)) \Delta @ ?T_0 \ominus (map \text{snd } \Delta)) \text{\$}\vdash \Phi$   
**using** *measure-deduction.simps(2)* **by** *blast*  
**let**  $?Psi' = \lambda \Psi \Delta$   
**let**  $?T_1 = map (uncurry (\rightarrow)) ?Psi' @ \Gamma \ominus (map \text{snd } ?Psi')$   
**let**  $?Delta' = \lambda \Psi \Delta$   
**have**  $(map (uncurry (\rightarrow)) ?Delta' @ ?T_1 \ominus (map \text{snd } ?Delta')) \text{\$}\vdash \Phi$   
 $map (uncurry (\sqcup)) \Psi \preceq map (uncurry (\sqcup)) ?Delta'$   
**using**  $\Psi(1) \Delta(1) \Delta(3)$   
 $xwitness-ywitness-measure-deduction-intro$   
 $Y-witness-left-stronger-theory$   
**by** *auto*  
**hence**  $?T_1 \text{\$}\vdash (\varphi \# \Phi)$   
**using**  $\Psi(1) \Psi(2) \Delta(1)$   
 $Y-witness-msub$  *measure-deduction.simps(2)*  
 $stronger-theory-deduction-monotonic$   
**by** *blast*  
**thus** *?thesis*  
**using**  $\Psi(1) \Delta(1) \Delta(2)$   
 $X-witness-msub$   
 $X-witness-right-stronger-theory$   
 $measure-deduction.simps(2)$   
 $stronger-theory-deduction-monotonic$   
**by** *blast*  
**qed**

**lemma** (in *classical-logic*) *measure-cons-remove1*:  
**assumes**  $\varphi \in set \Phi$   
**shows**  $\Gamma \text{\$}\vdash \Phi = \Gamma \text{\$}\vdash (\varphi \# (remove1 \varphi \Phi))$   
**proof** –  
**from**  $\langle \varphi \in set \Phi \rangle$   
**have**  $\forall \Gamma. \Gamma \text{\$}\vdash \Phi = \Gamma \text{\$}\vdash (\varphi \# (remove1 \varphi \Phi))$   
**proof** (*induct*  $\Phi$ )  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons*  $\chi \Phi$ )  
{  
**fix**  $\Gamma$

```

have  $\Gamma \Vdash (\chi \# \Phi) = \Gamma \Vdash (\varphi \# (\text{remove1 } \varphi (\chi \# \Phi)))$ 
proof (cases  $\chi = \varphi$ )
  case True
    then show ?thesis by simp
  next
    case False
      hence  $\varphi \in \text{set } \Phi$ 
        using Cons.prem by simp
      with Cons.hyps have  $\Gamma \Vdash (\chi \# \Phi) = \Gamma \Vdash (\chi \# \varphi \# (\text{remove1 } \varphi \Phi))$ 
        by fastforce
      hence  $\Gamma \Vdash (\chi \# \Phi) = \Gamma \Vdash (\varphi \# \chi \# (\text{remove1 } \varphi \Phi))$ 
        using measure-cons-cons-right-permute by blast
      then show ?thesis using  $\langle \chi \neq \varphi \rangle$  by simp
    qed
  }
  then show ?case by blast
qed
thus ?thesis using assms by blast
qed

lemma (in classical-logic) witness-stronger-theory:
  assumes  $\text{mset } (\text{map } \text{snd } \Psi) \subseteq\# \text{mset } \Gamma$ 
  shows  $(\text{map } (\text{uncurry } (\rightarrow))) \Psi @ \Gamma \ominus (\text{map } \text{snd } \Psi) \preceq \Gamma$ 
proof -
  have  $\forall \Gamma. \text{mset } (\text{map } \text{snd } \Psi) \subseteq\# \text{mset } \Gamma \longrightarrow (\text{map } (\text{uncurry } (\rightarrow))) \Psi @ \Gamma \ominus$ 
   $(\text{map } \text{snd } \Psi) \preceq \Gamma$ 
  proof (induct  $\Psi$ )
    case Nil
      then show ?case by simp
    next
      case (Cons  $\psi \Psi$ )
        let  $? \gamma = \text{snd } \psi$ 
        {
          fix  $\Gamma$ 
          assume  $\text{mset } (\text{map } \text{snd } (\psi \# \Psi)) \subseteq\# \text{mset } \Gamma$ 
          hence  $\text{mset } (\text{map } \text{snd } \Psi) \subseteq\# \text{mset } (\text{remove1 } (\text{snd } \psi) \Gamma)$ 
            by (simp add: insert-subset-eq-iff)
          with Cons have
             $(\text{map } (\text{uncurry } (\rightarrow))) \Psi @ (\text{remove1 } (\text{snd } \psi) \Gamma) \ominus (\text{map } \text{snd } \Psi) \preceq (\text{remove1 } ? \gamma \Gamma)$ 
            by blast
          hence  $(\text{map } (\text{uncurry } (\rightarrow))) \Psi @ \Gamma \ominus (\text{map } \text{snd } (\psi \# \Psi)) \preceq (\text{remove1 } ? \gamma \Gamma)$ 
            by (simp add: stronger-theory-relation-alt-def)
          moreover
          have  $(\text{uncurry } (\rightarrow)) = (\lambda \psi. \text{fst } \psi \rightarrow \text{snd } \psi)$ 
            by fastforce
          hence  $\vdash ? \gamma \rightarrow \text{uncurry } (\rightarrow) \psi$ 
            using axiom-k by simp
          ultimately have

```

```

    (map (uncurry (→)) (ψ # Ψ) @ Γ ⊖ (map snd (ψ # Ψ))) ≤ (?γ # (remove1
?γ Γ))
    using stronger-theory-left-right-cons by auto
    hence (map (uncurry (→)) (ψ # Ψ) @ Γ ⊖ (map snd (ψ # Ψ))) ≤ Γ
    using stronger-theory-relation-alt-def
      ⟨mset (map snd (ψ # Ψ)) ⊆# mset Γ⟩
      mset-subset-eqD
    by fastforce
  }
  then show ?case by blast
qed
thus ?thesis using assms by blast
qed

```

lemma (in classical-logic) measure-msub-weaken:

```

  assumes mset Ψ ⊆# mset Φ
    and Γ $⊢ Φ
  shows Γ $⊢ Ψ
proof -
  have ∀Ψ Γ. mset Ψ ⊆# mset Φ → Γ $⊢ Φ → Γ $⊢ Ψ
proof (induct Φ)
  case Nil
  then show ?case by simp
next
  case (Cons φ Φ)
  {
    fix Ψ Γ
    assume mset Ψ ⊆# mset (φ # Φ)
      Γ $⊢ (φ # Φ)
    hence Γ $⊢ Φ
      using measure-deduction.simps(2)
        measure-stronger-theory-left-monotonic
        witness-stronger-theory
      by blast
    have Γ $⊢ Ψ
  proof (cases φ ∈ set Ψ)
    case True
    hence mset (remove1 φ Ψ) ⊆# mset Φ
      using ⟨mset Ψ ⊆# mset (φ # Φ)⟩
        subset-eq-diff-conv
      by force
    hence ∀Γ. Γ $⊢ Φ → Γ $⊢ (remove1 φ Ψ)
      using Cons by blast
    hence Γ $⊢ (φ # (remove1 φ Ψ))
      using ⟨Γ $⊢ (φ # Φ)⟩ by fastforce
    then show ?thesis
      using ⟨φ ∈ set Ψ⟩
        measure-cons-remove1
      by blast
  }

```



```

next
  case False
  have  $mset \Psi \subseteq\# mset \Phi + add\text{-}mset \varphi (mset [])$ 
    using  $\langle mset \Psi \subseteq\# mset (\varphi \# \Phi) \rangle$  by auto
  hence  $mset \Psi \subseteq\# mset \Phi$ 
    by (metis (no-types) False
        diff-single-trivial
        in-multiset-in-set mset.simps(1)
        subset-eq-diff-conv)
  then show ?thesis
    using  $\langle \Gamma \ \$\vdash \Phi \rangle$  Cons
    by blast
qed
}
then show ?case by blast
qed
with assms show ?thesis by blast
qed

lemma (in classical-logic) measure-stronger-theory-right-antitonic:
  assumes  $\Psi \preceq \Phi$ 
  and  $\Gamma \ \$\vdash \Phi$ 
  shows  $\Gamma \ \$\vdash \Psi$ 
proof -
  have  $\forall \Psi \Gamma. \Psi \preceq \Phi \longrightarrow \Gamma \ \$\vdash \Phi \longrightarrow \Gamma \ \$\vdash \Psi$ 
  proof (induct  $\Phi$ )
  case Nil
  then show ?case
    using measure-deduction.simps(1)
    stronger-theory-empty-list-intro
    by blast
  next
  case (Cons  $\varphi \Phi$ )
  {
  fix  $\Psi \Gamma$ 
  assume  $\Gamma \ \$\vdash (\varphi \# \Phi)$ 
   $\Psi \preceq (\varphi \# \Phi)$ 
  from this obtain  $\Sigma$  where
     $\Sigma: map \text{snd } \Sigma = \Psi$ 
     $mset (map \text{fst } \Sigma) \subseteq\# mset (\varphi \# \Phi)$ 
     $\forall (\varphi, \psi) \in set \Sigma. \vdash \varphi \rightarrow \psi$ 
  unfolding stronger-theory-relation-def
  by auto
  hence  $\Gamma \ \$\vdash \Psi$ 
  proof (cases  $\varphi \in set (map \text{fst } \Sigma)$ )
  case True
  from this obtain  $\psi$  where  $(\varphi, \psi) \in set \Sigma$ 
  by (induct  $\Sigma$ , simp, fastforce)
  hence A:  $mset (map \text{snd } (remove1 (\varphi, \psi) \Sigma)) = mset (remove1 \psi \Psi)$ 

```

```

and  $B: mset (map fst (remove1 (\varphi, \psi) \Sigma)) \subseteq\# mset \Phi$ 
using  $\Sigma$  remove1-pairs-list-projections-snd
      remove1-pairs-list-projections-fst
      subset-eq-diff-conv
by fastforce+
have  $\forall (\varphi, \psi) \in set (remove1 (\varphi, \psi) \Sigma). \vdash \varphi \rightarrow \psi$ 
using  $\Sigma(3)$  by fastforce+
hence  $(remove1 \psi \Psi) \preceq \Phi$ 
unfolding stronger-theory-relation-alt-def using  $A B$  by blast
moreover
from  $\langle \Gamma \ \$\vdash (\varphi \# \Phi) \rangle$  obtain  $\Delta$  where
   $\Delta: mset (map snd \Delta) \subseteq\# mset \Gamma$ 
   $map (uncurry (\sqcup)) \Delta \vdash \varphi$ 
   $(map (uncurry (\rightarrow)) \Delta @ \Gamma \ominus (map snd \Delta)) \ \$\vdash \Phi$ 
by auto
ultimately have  $(map (uncurry (\rightarrow)) \Delta @ \Gamma \ominus (map snd \Delta)) \ \$\vdash remove1$ 
 $\psi \Psi$ 
using Cons by blast
moreover have  $map (uncurry (\sqcup)) \Delta \vdash \psi$ 
using  $\Delta(2)$   $\Sigma(3)$   $\langle (\varphi, \psi) \in set \Sigma \rangle$ 
      list-deduction-weaken
      list-deduction-modus-ponens
by blast
ultimately have  $\langle \Gamma \ \$\vdash (\psi \# (remove1 \psi \Psi)) \rangle$ 
using  $\Delta(1)$  by auto
moreover from  $\langle (\varphi, \psi) \in set \Sigma \rangle \Sigma(1)$  have  $\psi \in set \Psi$ 
by force
hence  $mset \Psi \subseteq\# mset (\psi \# (remove1 \psi \Psi))$ 
by auto
ultimately show ?thesis using measure-msub-weaken by blast
next
case False
hence  $mset (map fst \Sigma) \subseteq\# mset \Phi$ 
using  $\Sigma(2)$ 
by (simp,
      metis add-mset-add-single
      diff-single-trivial
      mset-map set-mset-mset
      subset-eq-diff-conv)
hence  $\Psi \preceq \Phi$ 
using  $\Sigma(1)$   $\Sigma(3)$ 
unfolding stronger-theory-relation-def
by auto
moreover from  $\langle \Gamma \ \$\vdash (\varphi \# \Phi) \rangle$  have  $\Gamma \ \$\vdash \Phi$ 
using measure-deduction.simps(2)
      measure-stronger-theory-left-monotonic
      witness-stronger-theory
by blast
ultimately show ?thesis using Cons by blast

```

```

    qed
  }
  then show ?case by blast
qed
thus ?thesis using assms by blast
qed

```

```

lemma (in classical-logic) measure-witness-right-split:
  assumes mset (map snd  $\Psi$ )  $\subseteq\#$  mset  $\Phi$ 
  shows  $\Gamma \ \$\vdash$  (map (uncurry ( $\sqcup$ ))  $\Psi$  @ map (uncurry ( $\rightarrow$ ))  $\Psi$  @  $\Phi$   $\ominus$  (map snd  $\Psi$ )) =  $\Gamma \ \$\vdash$   $\Phi$ 
proof -
  have  $\forall \Gamma \Phi. \text{mset (map snd } \Psi) \subseteq\# \text{mset } \Phi \longrightarrow$ 
     $\Gamma \ \$\vdash \Phi = \Gamma \ \$\vdash$  (map (uncurry ( $\sqcup$ ))  $\Psi$  @ map (uncurry ( $\rightarrow$ ))  $\Psi$  @  $\Phi$   $\ominus$  (map
  snd  $\Psi$ ))
  proof (induct  $\Psi$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\psi \Psi$ )
    {
      fix  $\Gamma \Phi$ 
      let ? $\chi$  = fst  $\psi$ 
      let ? $\varphi$  = snd  $\psi$ 
      let ? $\Phi'$  = map (uncurry ( $\sqcup$ )) ( $\psi \# \Psi$ ) @
        map (uncurry ( $\rightarrow$ )) ( $\psi \# \Psi$ ) @
         $\Phi$   $\ominus$  map snd ( $\psi \# \Psi$ )
      let ? $\Phi_0$  = map (uncurry ( $\sqcup$ ))  $\Psi$  @
        map (uncurry ( $\rightarrow$ ))  $\Psi$  @
        (remove1 ? $\varphi$   $\Phi$ )  $\ominus$  map snd  $\Psi$ 
      assume mset (map snd ( $\psi \# \Psi$ ))  $\subseteq\#$  mset  $\Phi$ 
      hence mset (map snd  $\Psi$ )  $\subseteq\#$  mset (remove1 ? $\varphi$   $\Phi$ )
        mset (? $\varphi \#$  remove1 ? $\varphi$   $\Phi$ ) = mset  $\Phi$ 
      by (simp add: insert-subset-eq-iff)+
      hence  $\Gamma \ \$\vdash \Phi = \Gamma \ \$\vdash$  (? $\varphi \#$  remove1 ? $\varphi$   $\Phi$ )
         $\forall \Gamma. \Gamma \ \$\vdash$  (remove1 ? $\varphi$   $\Phi$ ) =  $\Gamma \ \$\vdash$  ? $\Phi_0$ 
      by (metis list.set-intros(1) measure-cons-remove1 set-mset-mset,
        metis Cons.hyps)
    }
  moreover
  have (uncurry ( $\sqcup$ )) = ( $\lambda \psi. \text{fst } \psi \sqcup \text{snd } \psi$ )
    (uncurry ( $\rightarrow$ )) = ( $\lambda \psi. \text{fst } \psi \rightarrow \text{snd } \psi$ )
  by fastforce+
  hence mset ? $\Phi'$   $\subseteq\#$  mset (? $\chi \sqcup$  ? $\varphi \#$  ? $\chi \rightarrow$  ? $\varphi \#$  ? $\Phi_0$ )
    mset (? $\chi \sqcup$  ? $\varphi \#$  ? $\chi \rightarrow$  ? $\varphi \#$  ? $\Phi_0$ )  $\subseteq\#$  mset ? $\Phi'$ 
    (is mset ? $X \subseteq\#$  mset ? $Y$ )
  by fastforce+
  hence  $\Gamma \ \$\vdash$  ? $\Phi' = \Gamma \ \$\vdash$  (? $\varphi \#$  ? $\Phi_0$ )
  using measure-formula-right-split
    measure-msub-weaken

```

```

    by blast
  ultimately have  $\Gamma \Vdash \Phi = \Gamma \Vdash ?\Phi'$ 
    by fastforce
}
then show ?case by blast
qed
with assms show ?thesis by blast
qed

primrec (in classical-logic)
  submerge-witness :: ('a × 'a) list ⇒ ('a × 'a) list ⇒ ('a × 'a) list (⟨ $\mathfrak{E}$ ⟩)
where
   $\mathfrak{E} \Sigma [] = \text{map } (\lambda \sigma. (\perp, (\text{uncurry } (\sqcup)) \sigma)) \Sigma$ 
|  $\mathfrak{E} \Sigma (\delta \# \Delta) =$ 
  (case find  $(\lambda \sigma. (\text{uncurry } (\rightarrow)) \sigma = \text{snd } \delta) \Sigma$  of
    None ⇒  $\mathfrak{E} \Sigma \Delta$ 
  | Some  $\sigma \Rightarrow (\text{fst } \sigma, (\text{fst } \delta \sqcap \text{fst } \sigma) \sqcup \text{snd } \sigma) \# (\mathfrak{E} (\text{remove1 } \sigma \Sigma) \Delta))$ )

lemma (in classical-logic) submerge-witness-stronger-theory-left:
  map (uncurry  $(\sqcup)$ )  $\Sigma \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{E} \Sigma \Delta)$ 
proof -
  have  $\forall \Sigma. \text{map } (\text{uncurry } (\sqcup)) \Sigma \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{E} \Sigma \Delta)$ 
  proof (induct  $\Delta$ )
    case Nil
    {
      fix  $\Sigma$ 
      {
        fix  $\varphi$ 
        have  $\vdash (\perp \sqcup \varphi) \rightarrow \varphi$ 
          unfolding disjunction-def
          using ex-falso-quodlibet modus-ponens excluded-middle-elimination by blast
      }
      note tautology = this
      have map (uncurry  $(\sqcup)$ )  $\Sigma \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{E} \Sigma [])$ 
        by (induct  $\Sigma$ ,
            simp,
            simp add: stronger-theory-left-right-cons tautology)
    }
  then show ?case by auto
next
  case (Cons  $\delta \Delta$ )
  {
    fix  $\Sigma$ 
    have map (uncurry  $(\sqcup)$ )  $\Sigma \preceq \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{E} \Sigma (\delta \# \Delta))$ 
    proof (cases find  $(\lambda \sigma. (\text{uncurry } (\rightarrow)) \sigma = \text{snd } \delta) \Sigma = \text{None}$ )
      case True
      then show ?thesis using Cons by simp
    next
      case False

```

```

from this obtain  $\sigma$  where
 $\sigma$ : find ( $\lambda\sigma$ . uncurry ( $\rightarrow$ )  $\sigma$  = snd  $\delta$ )  $\Sigma$  = Some  $\sigma$ 
  uncurry ( $\rightarrow$ )  $\sigma$  = snd  $\delta$ 
   $\sigma \in \text{set } \Sigma$ 
using find-Some-predicate find-Some-set-membership
by fastforce
{
  fix  $\alpha \beta \gamma$ 
  have  $\vdash (\alpha \sqcup (\gamma \sqcap \alpha) \sqcup \beta) \rightarrow (\alpha \sqcup \beta)$ 
  proof –
    let  $?\varphi = (\langle\alpha\rangle \sqcup (\langle\gamma\rangle \sqcap \langle\alpha\rangle) \sqcup \langle\beta\rangle) \rightarrow (\langle\alpha\rangle \sqcup \langle\beta\rangle)$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\langle ?\varphi \rangle)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
}
note tautology = this
let  $?\alpha = \text{fst } \sigma$ 
let  $?\beta = \text{snd } \sigma$ 
let  $?\gamma = \text{fst } \delta$ 
have (uncurry ( $\sqcup$ )) = ( $\lambda \sigma$ .  $\text{fst } \sigma \sqcup \text{snd } \sigma$ ) by fastforce
hence (uncurry ( $\sqcup$ ))  $\sigma$  =  $?\alpha \sqcup ?\beta$  by simp
hence  $A$ :  $\vdash (?\alpha \sqcup (?\gamma \sqcap ?\alpha) \sqcup ?\beta) \rightarrow (\text{uncurry } (\sqcup)) \sigma$  using tautology by
simp
moreover
have map (uncurry ( $\sqcup$ )) (remove1  $\sigma \Sigma$ )
   $\preceq$  map (uncurry ( $\sqcup$ )) ( $\mathfrak{E}$  (remove1  $\sigma \Sigma$ )  $\Delta$ )
  using Cons by simp
ultimately have  $A$ :
  map (uncurry ( $\sqcup$ )) ( $\sigma \#$  (remove1  $\sigma \Sigma$ ))
   $\preceq$  ( $?\alpha \sqcup (?\gamma \sqcap ?\alpha) \sqcup ?\beta \#$  map (uncurry ( $\sqcup$ )) ( $\mathfrak{E}$  (remove1  $\sigma \Sigma$ )  $\Delta$ ))
  using stronger-theory-left-right-cons by fastforce
from  $\sigma(\beta)$  have mset  $\Sigma$  = mset ( $\sigma \#$  (remove1  $\sigma \Sigma$ ))
  by simp
  hence mset (map (uncurry ( $\sqcup$ ))  $\Sigma$ ) = mset (map (uncurry ( $\sqcup$ )) ( $\sigma \#$ 
(remove1  $\sigma \Sigma$ )))
  by (metis mset-map)
hence  $B$ : map (uncurry ( $\sqcup$ ))  $\Sigma \preceq$  map (uncurry ( $\sqcup$ )) ( $\sigma \#$  (remove1  $\sigma \Sigma$ ))
  by (simp add: msub-stronger-theory-intro)
have (  $\text{fst } \sigma$ 
   $\sqcup$  ( $\text{fst } \delta \sqcap \text{fst } \sigma$ )
   $\sqcup$   $\text{snd } \sigma \#$  map ( $\lambda(x, y). x \sqcup y$ ) ( $\mathfrak{E}$  (remove1  $\sigma \Sigma$ )  $\Delta$ ))  $\succeq$  map ( $\lambda(x,$ 
 $y). x \sqcup y$ )  $\Sigma$ 
  by (metis
    (no-types, lifting)
     $A B$ 
    stronger-theory-transitive
    uncurry-def)
thus ?thesis using  $A B \sigma$  by simp

```

```

    qed
  }
  then show ?case by auto
  qed
  thus ?thesis by blast
  qed

```

```

lemma (in classical-logic) submerge-witness-msub:
  mset (map snd (⊗ Σ Δ)) ⊆# mset (map (uncurry (⊔)) (⋈ Σ Δ))
proof –
  have ∀ Σ. mset (map snd (⊗ Σ Δ)) ⊆# mset (map (uncurry (⊔)) (⋈ Σ Δ))
proof (induct Δ)
  case Nil
  {
    fix Σ
    have mset (map snd (⊗ Σ [])) ⊆#
      mset (map (uncurry (⊔)) (⋈ Σ []))
      by (induct Σ, simp+)
  }
  then show ?case by blast
next
  case (Cons δ Δ)
  {
    fix Σ
    have mset (map snd (⊗ Σ (δ # Δ))) ⊆#
      mset (map (uncurry (⊔)) (⋈ Σ (δ # Δ)))
      using Cons
      by (cases find (λ σ. (uncurry (→)) σ = snd δ) Σ = None,
        simp,
        meson diff-subset-eq-self
          insert-subset-eq-iff
          mset-subset-eq-add-mset-cancel
          subset-mset.dual-order.trans,
        fastforce)
  }
  then show ?case by blast
  qed
  thus ?thesis by blast
  qed

```

```

lemma (in classical-logic) submerge-witness-stronger-theory-right:
  map (uncurry (⊔)) Δ
  ≼ (map (uncurry (→)) (⊗ Σ Δ) @ map (uncurry (⊔)) (⋈ Σ Δ) ⊖ map snd (⊗ Σ
  Δ))
proof –
  have ∀ Σ. map (uncurry (⊔)) Δ
    ≼ (map (uncurry (→)) (⊗ Σ Δ) @ map (uncurry (⊔)) (⋈ Σ Δ) ⊖ map
  snd (⊗ Σ Δ))
proof(induct Δ)

```

```

case Nil
then show ?case by simp
next
case (Cons δ Δ)
{
  fix Σ
  have map (uncurry (⊔)) (δ # Δ) ≤
    ( map (uncurry (→)) (⊔ Σ (δ # Δ))
      @ map (uncurry (⊔)) (⋈ Σ (δ # Δ))
        ⊖ map snd (⊔ Σ (δ # Δ)))
  proof (cases find (λ σ. (uncurry (→)) σ = snd δ) Σ = None)
  case True
  from Cons obtain Φ where Φ:
    map snd Φ = map (uncurry (⊔)) Δ
    mset (map fst Φ) ⊆#
      mset (map (uncurry (→)) (⊔ Σ Δ))
        @ map (uncurry (⊔)) (⋈ Σ Δ) ⊖ map snd (⊔ Σ Δ)
    ∀(γ, σ) ∈ set Φ. ⊢ γ → σ
  unfolding stronger-theory-relation-def
  by fastforce
  let ?Φ' = (uncurry (⊔)) δ, (uncurry (⊔)) δ # Φ
  have map snd ?Φ' = map (uncurry (⊔)) (δ # Δ) using Φ(1) by simp
  moreover
  from Φ(2) have A:
    image-mset fst (mset Φ)
    ⊆# {#x → y. (x, y) ∈# mset (⊔ Σ Δ)#}
      + ({#x ⊔ y. (x, y) ∈# mset (⋈ Σ Δ)#} - image-mset snd (mset (⊔ Σ
Δ)))
  by simp
  have image-mset snd (mset (⊔ Σ Δ)) ⊆# {#x ⊔ y. (x, y) ∈# mset (⋈ Σ
Δ)#}
  using submerge-witness-msub by force
  then have B: {#case δ of (x, xa) ⇒ x ⊔ xa#}
    ⊆# add-mset (case δ of (x, xa) ⇒ x ⊔ xa)
      {#x ⊔ y. (x, y) ∈# mset (⋈ Σ Δ)#} - image-mset snd
(mset (⊔ Σ Δ))
  by (metis add-mset-add-single subset-mset.le-add-diff)
  have add-mset (case δ of (x, xa) ⇒ x ⊔ xa) {#x ⊔ y. (x, y) ∈# mset (⋈
Σ Δ)#}
    - image-mset snd (mset (⊔ Σ Δ)) - {#case δ of (x, xa) ⇒ x ⊔ xa#}
    = {#x ⊔ y. (x, y) ∈# mset (⋈ Σ Δ)#} - image-mset snd (mset (⊔ Σ
Δ))
  by force
  then have add-mset (case δ of (x, xa) ⇒ x ⊔ xa) (image-mset fst (mset
Φ))
    - (add-mset (case δ of (x, xa) ⇒ x ⊔ xa) {#x ⊔ y. (x, y) ∈# mset
(⋈ Σ Δ)#}
      - image-mset snd (mset (⊔ Σ Δ)))
    ⊆# {#x → y. (x, y) ∈# mset (⊔ Σ Δ)#}

```

```

using  $A B$  by (metis (no-types) add-mset-add-single
                    subset-eq-diff-conv
                    subset-mset.diff-diff-right)
hence add-mset (case  $\delta$  of ( $x, xa$ )  $\Rightarrow x \sqcup xa$ ) (image-mset fst (mset  $\Phi$ ))
   $\subseteq\# \{ \#x \rightarrow y. (x, y) \in\# \text{mset } (\mathfrak{E} \Sigma \Delta)\# \}$ 
  + (add-mset (case  $\delta$  of ( $x, xa$ )  $\Rightarrow x \sqcup xa$ )  $\{ \#x \sqcup y. (x, y) \in\# \text{mset } (\mathfrak{J}$ 
 $\Sigma \Delta)\# \}$ 
    - image-mset snd (mset ( $\mathfrak{E} \Sigma \Delta$ )))
using subset-eq-diff-conv by blast
hence
  mset (map fst  $?\Phi'$ )  $\subseteq\#$ 
    mset (map (uncurry ( $\rightarrow$ )) ( $\mathfrak{E} \Sigma (\delta \# \Delta)$ )
      @ map (uncurry ( $\sqcup$ )) ( $\mathfrak{J} \Sigma (\delta \# \Delta)$ )
       $\ominus$  map snd ( $\mathfrak{E} \Sigma (\delta \# \Delta)$ ))
using True  $\Phi(2)$ 
by simp
moreover have  $\forall (\gamma, \sigma) \in \text{set } ?\Phi'. \vdash \gamma \rightarrow \sigma$ 
using  $\Phi(3)$  trivial-implication by auto
ultimately show  $?\text{thesis}$ 
unfolding stronger-theory-relation-def
by blast
next
case False
from this obtain  $\sigma$  where
   $\sigma: \text{find } (\lambda\sigma. \text{uncurry } (\rightarrow) \sigma = \text{snd } \delta) \Sigma = \text{Some } \sigma$ 
  uncurry ( $\rightarrow$ )  $\sigma = \text{snd } \delta$ 
using find-Some-predicate
by fastforce
moreover from Cons have
  map (uncurry ( $\sqcup$ ))  $\Delta \preceq$ 
  (map (uncurry ( $\rightarrow$ )) ( $\mathfrak{E} (\text{remove1 } \sigma \Sigma) \Delta$ ) @
    remove1 ( $(\text{fst } \delta \sqcap \text{fst } \sigma) \sqcup \text{snd } \sigma$ )
    ( $(\text{fst } \delta \sqcap \text{fst } \sigma) \sqcup \text{snd } \sigma \# \text{map } (\text{uncurry } (\sqcup)) (\mathfrak{J} (\text{remove1 } \sigma \Sigma) \Delta)$ )
     $\ominus$  map snd ( $\mathfrak{E} (\text{remove1 } \sigma \Sigma) \Delta$ ))
unfolding stronger-theory-relation-alt-def
by simp
moreover
{
  fix  $\alpha \beta \gamma$ 
  have  $\vdash (\alpha \rightarrow ((\gamma \sqcap \alpha) \sqcup \beta)) \rightarrow (\gamma \sqcup (\alpha \rightarrow \beta))$ 
  proof -
    let  $?\varphi = (\langle \alpha \rangle \rightarrow ((\langle \gamma \rangle \sqcap \langle \alpha \rangle) \sqcup \langle \beta \rangle)) \rightarrow (\langle \gamma \rangle \sqcup (\langle \alpha \rangle \rightarrow \langle \beta \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\langle ?\varphi \rangle)$  using propositional-semantic by blast
    thus  $?\text{thesis}$  by simp
  qed
}
note tautology = this
let  $?\alpha = \text{fst } \sigma$ 

```



```

    let ?β = snd σ
    let ?γ = fst δ
    have (λ δ. uncurry (⊔) δ) = (λ δ. fst δ ⊔ snd δ)
      (λ σ. uncurry (→) σ) = (λ σ. fst σ → snd σ) by fastforce+
    hence (uncurry (⊔) δ) = (?γ ⊔ (?α → ?β)) using σ(2) by simp
    hence ⊢ (?α → ((?γ ⊓ ?α) ⊔ ?β)) → (uncurry (⊔) δ) using tautology by
  auto
    ultimately show ?thesis
      using stronger-theory-left-right-cons
      by fastforce
    qed
  }
  then show ?case by auto
  qed
  thus ?thesis by simp
  qed

```

**lemma** (in *classical-logic*) *merge-witness-cons-measure-deduction*:

```

  assumes map (uncurry (⊔)) Σ ⊢ φ
    and mset (map snd Δ) ⊆# mset (map (uncurry (→)) Σ @ Γ ⊖ map snd Σ)
    and map (uncurry (⊔)) Δ $⊢ Φ
    shows map (uncurry (⊔)) (⋈ Σ Δ) $⊢ (φ # Φ)
  proof -
    let ?Σ' = ⋈ Σ Δ
    let ?Γ = map (uncurry (→)) ?Σ' @ map (uncurry (⊔)) (⋈ Σ Δ) ⊖ map snd ?Σ'
    have ?Γ $⊢ Φ
      using assms(3)
      submerge-witness-stronger-theory-right
      measure-stronger-theory-left-monotonic
      by blast
    moreover have map (uncurry (⊔)) ?Σ' ⊢ φ
      using assms(1)
      stronger-theory-deduction-monotonic
      submerge-witness-stronger-theory-left
      by blast
    ultimately show ?thesis
      using submerge-witness-msub
      by fastforce
  qed

```

**primrec** (in *classical-logic*)

*recover-witness-A* :: ('a × 'a) list ⇒ ('a × 'a) list ⇒ ('a × 'a) list (⟨℘⟩)

**where**

```

  ℘ Σ [] = Σ
| ℘ Σ (δ # Δ) =
  (case find (λ σ. snd σ = (uncurry (⊔)) δ) Σ of
   None ⇒ ℘ Σ Δ
  | Some σ ⇒ (fst σ ⊔ fst δ, snd δ) # (℘ (remove1 σ Σ) Δ))

```

**primrec** (in *classical-logic*)  
*recover-complement-A* :: ('a × 'a) list ⇒ ('a × 'a) list ⇒ ('a × 'a) list (⟨ $\mathfrak{P}^C$ ⟩)  
**where**  
 $\mathfrak{P}^C \Sigma [] = []$   
|  $\mathfrak{P}^C \Sigma (\delta \# \Delta) =$   
    (case find (λ σ. snd σ = (uncurry (⊔)) δ) Σ of  
      None ⇒ δ #  $\mathfrak{P}^C \Sigma \Delta$   
      Some σ ⇒  $\mathfrak{P}^C$  (remove1 σ Σ) Δ))

**primrec** (in *classical-logic*)  
*recover-witness-B* :: ('a × 'a) list ⇒ ('a × 'a) list ⇒ ('a × 'a) list (⟨ $\mathfrak{Q}$ ⟩)  
**where**  
 $\mathfrak{Q} \Sigma [] = []$   
|  $\mathfrak{Q} \Sigma (\delta \# \Delta) =$   
    (case find (λ σ. (snd σ) = (uncurry (⊔)) δ) Σ of  
      None ⇒ δ #  $\mathfrak{Q} \Sigma \Delta$   
      Some σ ⇒ (fst δ, (fst σ ⊔ fst δ) → snd δ) # ( $\mathfrak{Q}$  (remove1 σ Σ) Δ))

**lemma** (in *classical-logic*) *recover-witness-A-left-stronger-theory*:  
map (uncurry (⊔)) Σ ≲ map (uncurry (⊔)) ( $\mathfrak{P} \Sigma \Delta$ )

**proof** –

**have** ∀ Σ. map (uncurry (⊔)) Σ ≲ map (uncurry (⊔)) ( $\mathfrak{P} \Sigma \Delta$ )

**proof** (induct Δ)

case Nil

{

  fix Σ

**have** map (uncurry (⊔)) Σ ≲ map (uncurry (⊔)) ( $\mathfrak{P} \Sigma []$ )

**by**(induct Σ, simp+)

}

**then show** ?case by auto

**next**

case (Cons δ Δ)

{

  fix Σ

**have** map (uncurry (⊔)) Σ ≲ map (uncurry (⊔)) ( $\mathfrak{P} \Sigma (\delta \# \Delta)$ )

**proof** (cases find (λ σ. snd σ = uncurry (⊔) δ) Σ = None)

    case True

**then show** ?thesis using Cons by simp

**next**

    case False

**from this obtain** σ **where**

      σ: find (λσ. snd σ = uncurry (⊔) δ) Σ = Some σ

      snd σ = uncurry (⊔) δ

      σ ∈ set Σ

**using** find-Some-predicate

      find-Some-set-membership

**by** fastforce

**let** ?α = fst σ

**let** ?β = fst δ

```

let ?γ = snd δ
have uncurry (⊔) = (λδ. fst δ ⊔ snd δ) by fastforce
hence ⊢ ((?α ⊔ ?β) ⊔ ?γ) → uncurry (⊔) σ
  using σ(2) biconditional-def disjunction-associativity
  by auto
moreover
have map (uncurry (⊔)) (remove1 σ Σ)
  ≤ map (uncurry (⊔)) (⌘ (remove1 σ Σ) Δ)
  using Cons by simp
ultimately have map (uncurry (⊔)) (σ # (remove1 σ Σ))
  ≤ map (uncurry (⊔)) (⌘ Σ (δ # Δ))
  using σ(1)
  by (simp, metis stronger-theory-left-right-cons)
moreover
from σ(3) have mset Σ = mset (σ # (remove1 σ Σ))
  by simp
  hence mset (map (uncurry (⊔)) Σ) = mset (map (uncurry (⊔)) (σ #
(remove1 σ Σ)))
  by (metis mset-map)
  hence map (uncurry (⊔)) Σ ≤ map (uncurry (⊔)) (σ # (remove1 σ Σ))
  by (simp add: msub-stronger-theory-intro)
ultimately show ?thesis
  using stronger-theory-transitive by blast
qed
}
then show ?case by blast
qed
thus ?thesis by auto
qed

lemma (in classical-logic) recover-witness-A-mset-equiv:
  assumes mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ)
  shows mset (map snd (⌘ Σ Δ @ ⌘C Σ Δ)) = mset (map snd Δ)
proof -
  have ∀ Σ. mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ)
    → mset (map snd (⌘ Σ Δ @ ⌘C Σ Δ)) = mset (map snd Δ)
  proof (induct Δ)
    case Nil
    then show ?case by simp
  next
    case (Cons δ Δ)
    {
      fix Σ :: ('a × 'a) list
      assume *: mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) (δ # Δ))
      have mset (map snd (⌘ Σ (δ # Δ) @ ⌘C Σ (δ # Δ))) = mset (map snd (δ
# Δ))
      proof (cases find (λ σ. snd σ = uncurry (⊔) δ) Σ = None)
        case True
        hence uncurry (⊔) δ ∉ set (map snd Σ)

```

```

proof (induct  $\Sigma$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\sigma$   $\Sigma$ )
  then show ?case
    by (cases (uncurry ( $\sqcup$ ))  $\delta = \text{snd } \sigma$ , fastforce+)
  qed
  moreover have  $mset (\text{map } \text{snd } \Sigma) \subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) \Delta) +$ 
 $\{\#\text{uncurry } (\sqcup) \delta\#$ 
    using  $\star$  by fastforce
  ultimately have  $mset (\text{map } \text{snd } \Sigma) \subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) \Delta)$ 
    by (metis diff-single-trivial
      in-multiset-in-set
      subset-eq-diff-conv)
  then show ?thesis using Cons True by simp
next
  case False
  from this obtain  $\sigma$  where
     $\sigma$ : find ( $\lambda\sigma. \text{snd } \sigma = \text{uncurry } (\sqcup) \delta$ )  $\Sigma = \text{Some } \sigma$ 
     $\text{snd } \sigma = \text{uncurry } (\sqcup) \delta$ 
     $\sigma \in \text{set } \Sigma$ 
  using find-Some-predicate
    find-Some-set-membership
  by fastforce
  have  $A$ :  $mset (\text{map } \text{snd } \Sigma)$ 
     $\subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) \Delta) + \text{add-mset } (\text{uncurry } (\sqcup) \delta) (mset [])$ 
    using  $\star$  by auto
  have ( $\text{fst } \sigma, \text{uncurry } (\sqcup) \delta$ )  $\in\# mset \Sigma$ 
    by (metis (no-types)  $\sigma(2)$   $\sigma(3)$  prod.collapse set-mset-mset)
  then have  $B$ :  $mset (\text{map } \text{snd } (\text{remove1 } (\text{fst } \sigma, \text{uncurry } (\sqcup) \delta) \Sigma))$ 
     $= mset (\text{map } \text{snd } \Sigma) - \{\#\text{uncurry } (\sqcup) \delta\#$ 
    by (meson remove1-pairs-list-projections-snd)
  have ( $\text{fst } \sigma, \text{uncurry } (\sqcup) \delta$ )  $= \sigma$ 
    by (metis  $\sigma(2)$  prod.collapse)
  then have  $mset (\text{map } \text{snd } \Sigma) - \text{add-mset } (\text{uncurry } (\sqcup) \delta) (mset [])$ 
     $= mset (\text{map } \text{snd } (\text{remove1 } \sigma \Sigma))$ 
    using  $B$  by simp
  hence  $mset (\text{map } \text{snd } (\text{remove1 } \sigma \Sigma)) \subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) \Delta)$ 
    using  $A$  by (metis (no-types) subset-eq-diff-conv)
  with  $\sigma(1)$  Cons show ?thesis by simp
  qed
}
then show ?case by simp
qed
with assms show ?thesis by blast
qed

```

**lemma** (*in classical-logic*) *recover-witness-B-stronger-theory*:

```

assumes  $mset (map\ snd\ \Sigma) \subseteq\# mset (map (uncurry (\sqcup)) \Delta)$ 
shows  $(map (uncurry (\rightarrow)) \Sigma @ map (uncurry (\sqcup)) \Delta \ominus map\ snd\ \Sigma)$ 
 $\preceq map (uncurry (\sqcup)) (\Omega\ \Sigma\ \Delta)$ 
proof -
have  $\forall\ \Sigma. mset (map\ snd\ \Sigma) \subseteq\# mset (map (uncurry (\sqcup)) \Delta)$ 
 $\rightarrow (map (uncurry (\rightarrow)) \Sigma @ map (uncurry (\sqcup)) \Delta \ominus map\ snd\ \Sigma)$ 
 $\preceq map (uncurry (\sqcup)) (\Omega\ \Sigma\ \Delta)$ 
proof (induct  $\Delta$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\delta\ \Delta$ )
  {
    fix  $\Sigma :: ('a \times 'a)\ list$ 
    assume  $\star: mset (map\ snd\ \Sigma) \subseteq\# mset (map (uncurry (\sqcup)) (\delta \# \Delta))$ 
    have  $(map (uncurry (\rightarrow)) \Sigma @ map (uncurry (\sqcup)) (\delta \# \Delta) \ominus map\ snd\ \Sigma)$ 
 $\preceq map (uncurry (\sqcup)) (\Omega\ \Sigma\ (\delta \# \Delta))$ 
    proof (cases find  $(\lambda\ \sigma. snd\ \sigma = uncurry (\sqcup)\ \delta)\ \Sigma = None$ )
      case True
      hence  $uncurry (\sqcup)\ \delta \notin set (map\ snd\ \Sigma)$ 
      proof (induct  $\Sigma$ )
        case Nil
        then show ?case by simp
      next
        case (Cons  $\sigma\ \Sigma$ )
        then show ?case
          by (cases uncurry  $(\sqcup)\ \delta = snd\ \sigma, fastforce+$ )
      qed
    hence  $mset (map (uncurry (\rightarrow)) \Sigma @ (map (uncurry (\sqcup)) (\delta \# \Delta)) \ominus map$ 
 $snd\ \Sigma)$ 
 $= mset (uncurry (\sqcup)\ \delta \# map (uncurry (\rightarrow)) \Sigma$ 
 $@ map (uncurry (\sqcup)) \Delta \ominus map\ snd\ \Sigma)$ 
 $mset (map\ snd\ \Sigma) \subseteq\# mset (map (uncurry (\sqcup)) \Delta)$ 
    using  $\star$ 
    by (simp, simp,
      metis add-mset-add-single
      diff-single-trivial
      image-set
      mset-map
      set-mset-mset
      subset-eq-diff-conv)
    moreover from this have
 $(map (uncurry (\rightarrow)) \Sigma @ map (uncurry (\sqcup)) \Delta \ominus map\ snd\ \Sigma)$ 
 $\preceq map (uncurry (\sqcup)) (\Omega\ \Sigma\ \Delta)$ 
    using Cons
    by auto
    hence  $(uncurry (\sqcup)\ \delta \# map (uncurry (\rightarrow)) \Sigma @ map (uncurry (\sqcup)) \Delta \ominus$ 
 $map\ snd\ \Sigma)$ 
 $\preceq map (uncurry (\sqcup)) (\Omega\ \Sigma\ (\delta \# \Delta))$ 

```

```

    using True
    by (simp add: stronger-theory-left-right-cons trivial-implication)
  ultimately show ?thesis
    unfolding stronger-theory-relation-alt-def
    by simp
next
case False
let ?Γ = map (uncurry (→)) Σ @ (map (uncurry (⊔)) (δ # Δ)) ⊖ map snd
Σ
from False obtain σ where
  σ: find (λσ. snd σ = uncurry (⊔) δ) Σ = Some σ
  snd σ = uncurry (⊔) δ
  σ ∈ set Σ
using find-Some-predicate
  find-Some-set-membership
by fastforce
let ?Γ₀ = map (uncurry (→)) (remove1 σ Σ)
  @ (map (uncurry (⊔)) Δ) ⊖ map snd (remove1 σ Σ)
let ?α = fst σ
let ?β = fst δ
let ?γ = snd δ
have uncurry (⊔) = (λ σ. fst σ ⊔ snd σ)
  uncurry (→) = (λ σ. fst σ → snd σ)
by fastforce+
hence uncurry (→) σ = ?α → (?β ⊔ ?γ)
  using σ(2)
  by simp
from σ(3) have mset (σ # (remove1 σ Σ)) = mset Σ by simp
hence ♠: mset (map snd (σ # (remove1 σ Σ))) = mset (map snd Σ)
  mset (map (uncurry (→)) (σ # (remove1 σ Σ))) = mset (map
(uncurry (→)) Σ)
  by (metis mset-map)+
hence mset ?Γ = mset (map (uncurry (→)) (σ # (remove1 σ Σ)))
  @ (uncurry (⊔) δ # map (uncurry (⊔)) Δ)
  ⊖ map snd (σ # (remove1 σ Σ))
  by simp
hence ?Γ ≼ (?α → (?β ⊔ ?γ) # ?Γ₀)
  using σ(2) ⟨uncurry (→) σ = ?α → (?β ⊔ ?γ)⟩
  by (simp add: msub-stronger-theory-intro)
moreover have mset (map snd (remove1 σ Σ)) ⊆# mset (map (uncurry
(⊔)) Δ)
  using ♠(1)
  by (simp,
    metis (no-types, lifting)
      * σ(2)
      list.simps(9)
      mset.simps(2)
      mset-map
      uncurry-def)

```

```

      mset-subset-eq-add-mset-cancel)
with Cons have  $\heartsuit$ :  $?T_0 \preceq \text{map } (\text{uncurry } (\sqcup)) (\heartsuit \text{ (remove1 } \sigma \Sigma) \Delta)$  by
simp
{
  fix  $\alpha \beta \gamma$ 
  have  $\vdash (\beta \sqcup (\alpha \sqcup \beta) \rightarrow \gamma) \rightarrow (\alpha \rightarrow (\beta \sqcup \gamma))$ 
  proof -
    let  $?\varphi = (\langle \beta \rangle \sqcup (\langle \alpha \rangle \sqcup \langle \beta \rangle) \rightarrow \langle \gamma \rangle) \rightarrow (\langle \alpha \rangle \rightarrow (\langle \beta \rangle \sqcup \langle \gamma \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\heartsuit ?\varphi)$  using propositional-semantic by blast
    thus  $?thesis$  by simp
  qed
}
hence  $\vdash (?\beta \sqcup (?\alpha \sqcup ?\beta) \rightarrow ?\gamma) \rightarrow (?\alpha \rightarrow (?\beta \sqcup ?\gamma))$ 
  by simp
hence  $(?\alpha \rightarrow (?\beta \sqcup ?\gamma)) \# ?T_0 \preceq \text{map } (\text{uncurry } (\sqcup)) (\heartsuit \Sigma (\delta \# \Delta))$ 
  using  $\sigma(1) \heartsuit$ 
  by (simp, metis stronger-theory-left-right-cons)
ultimately show  $?thesis$ 
  using stronger-theory-transitive by blast
qed
}
then show  $?case$  by simp
qed
thus  $?thesis$  using assms by blast
qed

```

```

lemma (in classical-logic) recover-witness-B-mset-equiv:
  assumes  $mset (\text{map snd } \Sigma) \subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) \Delta)$ 
  shows  $mset (\text{map snd } (\heartsuit \Sigma \Delta))$ 
    =  $mset (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{P} \Sigma \Delta) @ \text{map snd } \Delta \ominus \text{map snd } (\mathfrak{P} \Sigma \Delta))$ 
proof -
  have  $\forall \Sigma. mset (\text{map snd } \Sigma) \subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) \Delta)$ 
     $\rightarrow mset (\text{map snd } (\heartsuit \Sigma \Delta)) = mset (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{P} \Sigma \Delta) @$ 
map snd  $(\mathfrak{P}^C \Sigma \Delta))$ 
  proof (induct  $\Delta$ )
    case Nil
    then show  $?case$  by simp
  next
  case (Cons  $\delta \Delta$ )
  {
    fix  $\Sigma :: ('a \times 'a) \text{ list}$ 
    assume  $*$ :  $mset (\text{map snd } \Sigma) \subseteq\# mset (\text{map } (\text{uncurry } (\sqcup)) (\delta \# \Delta))$ 
    have  $mset (\text{map snd } (\heartsuit \Sigma (\delta \# \Delta)))$ 
      =  $mset (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{P} \Sigma (\delta \# \Delta)) @ \text{map snd } (\mathfrak{P}^C \Sigma (\delta \# \Delta)))$ 
    proof (cases find  $(\lambda \sigma. \text{snd } \sigma = \text{uncurry } (\sqcup) \delta) \Sigma = \text{None}$ )
      case True
      hence  $\text{uncurry } (\sqcup) \delta \notin \text{set } (\text{map snd } \Sigma)$ 
      proof (induct  $\Sigma$ )

```

```

    case Nil
    then show ?case by simp
next
    case (Cons σ Σ)
    then show ?case
      by (cases (uncurry (⊔)) δ = snd σ, fastforce+)
qed
    moreover have mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ) +
{#uncurry (⊔) δ#}
      using * by force
    ultimately have mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ)
      by (metis diff-single-trivial in-multiset-in-set subset-eq-diff-conv)
    then show ?thesis using True Cons by simp
next
    case False
    from this obtain σ where
      σ: find (λσ. snd σ = uncurry (⊔) δ) Σ = Some σ
      snd σ = uncurry (⊔) δ
      σ ∈ set Σ
    using find-Some-predicate
      find-Some-set-membership
    by fastforce
    hence (fst σ, uncurry (⊔) δ) ∈# mset Σ
      by (metis (full-types) prod.collapse set-mset-mset)
    then have mset (map snd (remove1 (fst σ, uncurry (⊔) δ) Σ))
      = mset (map snd Σ) - {#uncurry (⊔) δ#}
      by (meson remove1-pairs-list-projections-snd)
    moreover have
      mset (map snd Σ)
    ⊆# mset (map (uncurry (⊔)) Δ) + add-mset (uncurry (⊔) δ) (mset [])
      using * by force
    ultimately have mset (map snd (remove1 σ Σ))
      ⊆# mset (map (uncurry (⊔)) Δ)
      by (metis (no-types) σ(2) mset.simps(1) prod.collapse subset-eq-diff-conv)
    with σ(1) Cons show ?thesis by simp
  qed
}
then show ?case by blast
qed
thus ?thesis
  using assms recover-witness-A-mset-equiv
  by (simp, metis add-diff-cancel-left')
qed

lemma (in classical-logic) recover-witness-B-right-stronger-theory:
  map (uncurry (→)) Δ ≲ map (uncurry (→)) (⊔ Σ Δ)
proof -
  have ∀ Σ. map (uncurry (→)) Δ ≲ map (uncurry (→)) (⊔ Σ Δ)
  proof (induct Δ)

```



```

    case Nil
  then show ?case by simp
next
case (Cons δ Δ)
{
  fix Σ
  have map (uncurry (→)) (δ # Δ) ≃ map (uncurry (→)) (Ω Σ (δ # Δ))
  proof (cases find (λ σ. snd σ = uncurry (⊔) δ) Σ = None)
    case True
    then show ?thesis
      using Cons
      by (simp add: stronger-theory-left-right-cons trivial-implication)
  next
  case False
  from this obtain σ where σ:
    find (λ σ. snd σ = uncurry (⊔) δ) Σ = Some σ
  by fastforce
  let ?α = fst δ
  let ?β = snd δ
  let ?γ = fst σ
  have uncurry (→) = (λ δ. fst δ → snd δ) by fastforce
  hence uncurry (→) δ = ?α → ?β by auto
  moreover have ⊢ (?α → (?γ ⊔ ?α) → ?β) → ?α → ?β
    unfolding disjunction-def
    using axiom-k axiom-s modus-ponens flip-implication
    by blast
  ultimately show ?thesis
    using Cons σ
    by (simp add: stronger-theory-left-right-cons)
  qed
}
then show ?case by simp
qed
thus ?thesis by simp
qed

```

**lemma** (in *classical-logic*) *recoverWitnesses-mset-equiv*:

```

  assumes mset (map snd Δ) ⊆# mset Γ
  and mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ)
  shows mset (Γ ⊖ map snd Δ)
    = mset ((map (uncurry (→)) (℘ Σ Δ) @ Γ ⊖ map snd (℘ Σ Δ)) ⊖ map
  snd (Ω Σ Δ))
  proof -
  have mset (Γ ⊖ map snd Δ) = mset (Γ ⊖ map snd (℘C Σ Δ) ⊖ map snd (℘
  Σ Δ))
  using assms(2) recover-witness-A-mset-equiv
  by (simp add: union-commute)
  moreover have ∀ Σ. mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ)
    → mset (Γ ⊖ map snd (℘C Σ Δ))

```

```

      = (mset ((map (uncurry (→)) (℘ Σ Δ) @ Γ) ⊖ map snd (Ω Σ
Δ)))
  using assms(1)
  proof (induct Δ)
  case Nil
  then show ?case by simp
next
case (Cons δ Δ)
from Cons.prems have snd δ ∈ set Γ
  using mset-subset-eqD by fastforce
from Cons.prems have ♡: mset (map snd Δ) ⊆# mset Γ
  using subset-mset.dual-order.trans
  by fastforce
{
  fix Σ :: ('a × 'a) list
  assume *: mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) (δ # Δ))
  have mset (Γ ⊖ map snd (℘C Σ (δ # Δ)))
    = mset ((map (uncurry (→)) (℘ Σ (δ # Δ)) @ Γ) ⊖ map snd (Ω Σ (δ #
Δ)))
  proof (cases find (λ σ. snd σ = uncurry (⊔) δ) Σ = None)
  case True
  hence uncurry (⊔) δ ∉ set (map snd Σ)
  proof (induct Σ)
  case Nil
  then show ?case by simp
  next
  case (Cons σ Σ)
  then show ?case
    by (cases (uncurry (⊔)) δ = snd σ, fastforce+)
  qed
  moreover have mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ) +
  {#uncurry (⊔) δ#}
    using * by auto
  ultimately have mset (map snd Σ) ⊆# mset (map (uncurry (⊔)) Δ)
  by (metis (full-types) diff-single-trivial in-multiset-in-set subset-eq-diff-conv)
  with Cons.hyps ♡ have mset (Γ ⊖ map snd (℘C Σ Δ))
    = mset ((map (uncurry (→)) (℘ Σ Δ) @ Γ) ⊖ map snd
(Ω Σ Δ))
    by simp
  thus ?thesis using True ⟨snd δ ∈ set Γ⟩ by simp
}
next
case False
from this obtain σ where σ:
  find (λσ. snd σ = uncurry (⊔) δ) Σ = Some σ
  snd σ = uncurry (⊔) δ
  σ ∈ set Σ
  using find-Some-predicate
  find-Some-set-membership
  by fastforce

```

$\Delta$ ) **with**  $\star$  **have**  $mset (map\ snd (remove1\ \sigma\ \Sigma)) \subseteq\# mset (map (uncurry (\sqcup)))$   
**by** (*simp*, *metis* (*no-types*, *lifting*))  
*add-mset-remove-trivial-eq*  
*image-mset-add-mset*  
*in-multiset-in-set*  
*mset-subset-eq-add-mset-cancel*  
**with** *Cons.hyps* **have**  $mset (\Gamma \ominus map\ snd (\mathfrak{P}^C (remove1\ \sigma\ \Sigma)\ \Delta))$   
 $= mset ((map (uncurry (\rightarrow))) (\mathfrak{P} (remove1\ \sigma\ \Sigma)\ \Delta) @ \Gamma)$   
 $\ominus map\ snd (\mathfrak{Q} (remove1\ \sigma\ \Sigma)\ \Delta))$   
**using**  $\heartsuit$  **by** *blast*  
**then show** *?thesis* **using**  $\sigma$  **by** *simp*  
**qed**  
**}**  
**then show** *?case* **by** *blast*  
**qed**  
**moreover have**  $image\ mset\ snd (mset (\mathfrak{P}^C\ \Sigma\ \Delta)) = mset (map\ snd\ \Delta \ominus map\ snd (\mathfrak{P}\ \Sigma\ \Delta))$   
**using** *assms*(2) *recover-witness-A-mset-equiv*  
**by** (*simp*, *metis* (*no-types*)) *diff-union-cancelL* *list-subtract-mset-homomorphism* *mset-map*  
**then have**  $mset\ \Gamma - (image\ mset\ snd (mset (\mathfrak{P}^C\ \Sigma\ \Delta)) + image\ mset\ snd (mset (\mathfrak{P}\ \Sigma\ \Delta)))$   
 $= \{\#x \rightarrow y. (x, y) \in\# mset (\mathfrak{P}\ \Sigma\ \Delta)\#\}$   
 $+ (mset\ \Gamma - image\ mset\ snd (mset (\mathfrak{P}\ \Sigma\ \Delta))) - image\ mset\ snd (mset (\mathfrak{Q}\ \Sigma\ \Delta))$   
**using** *calculation*  
*assms*(2)  
*recover-witness-A-mset-equiv*  
*recover-witness-B-mset-equiv*  
**by** *fastforce*  
**ultimately**  
**show** *?thesis*  
**using** *assms* *recover-witness-A-mset-equiv*  
**by** *simp*  
**qed**

**theorem** (*in classical-logic*) *measure-deduction-generalized-witness*:  
 $\Gamma\ \$\vdash (\Phi @ \Psi) = (\exists\ \Sigma. mset (map\ snd\ \Sigma) \subseteq\# mset\ \Gamma \wedge$   
 $map (uncurry (\sqcup))\ \Sigma\ \$\vdash\ \Phi \wedge$   
 $(map (uncurry (\rightarrow))\ \Sigma @ \Gamma \ominus (map\ snd\ \Sigma))\ \$\vdash\ \Psi)$

**proof** –  
**have**  $\forall\ \Gamma\ \Psi. \Gamma\ \$\vdash (\Phi @ \Psi) = (\exists\ \Sigma. mset (map\ snd\ \Sigma) \subseteq\# mset\ \Gamma \wedge$   
 $map (uncurry (\sqcup))\ \Sigma\ \$\vdash\ \Phi \wedge$   
 $(map (uncurry (\rightarrow))\ \Sigma @ \Gamma \ominus (map\ snd\ \Sigma))\ \$\vdash\ \Psi)$

**proof** (*induct*  $\Phi$ )  
**case** *Nil*  
**{**  
**fix**  $\Gamma\ \Psi$

**have**  $\Gamma \text{\$}\vdash (\square @ \Psi) = (\exists \Sigma. \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
 $\text{map} (\text{uncurry } (\sqcup)) \Sigma \text{\$}\vdash \square \wedge$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \text{\$}\vdash \Psi)$

**proof** (*rule iffI*)  
**assume**  $\Gamma \text{\$}\vdash (\square @ \Psi)$   
**moreover**  
**have**  $\Gamma \text{\$}\vdash (\square @ \Psi) = (\text{mset} (\text{map snd } \square) \subseteq\# \text{mset } \Gamma \wedge$   
 $\text{map} (\text{uncurry } (\sqcup)) \square \text{\$}\vdash \square \wedge$   
 $\text{map} (\text{uncurry } (\rightarrow)) \square @ \Gamma \ominus (\text{map snd } \square) \text{\$}\vdash \Psi)$

**by simp**  
**ultimately show**  $\exists \Sigma. \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
 $\text{map} (\text{uncurry } (\sqcup)) \Sigma \text{\$}\vdash \square \wedge$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \text{\$}\vdash \Psi$

**by metis**

**next**  
**assume**  $\exists \Sigma. \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
 $\text{map} (\text{uncurry } (\sqcup)) \Sigma \text{\$}\vdash \square \wedge$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \text{\$}\vdash \Psi$

**from this obtain**  $\Sigma$  **where**  
 $\Sigma: \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \text{\$}\vdash (\square @ \Psi)$

**by fastforce**  
**hence**  $(\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma) \preceq \Gamma$   
**using witness-stronger-theory by auto**

**with**  $\Sigma(2)$  **show**  $\Gamma \text{\$}\vdash (\square @ \Psi)$   
**using measure-stronger-theory-left-monotonic by blast**

**qed**

**}**

**then show ?case by blast**

**next**

**case** (*Cons*  $\varphi \Phi$ )

**{**

**fix**  $\Gamma \Psi$

**have**  $\Gamma \text{\$}\vdash ((\varphi \# \Phi) @ \Psi) = (\exists \Sigma. \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
 $\text{map} (\text{uncurry } (\sqcup)) \Sigma \text{\$}\vdash (\varphi \# \Phi) \wedge$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \text{\$}\vdash \Psi)$

**proof** (*rule iffI*)  
**assume**  $\Gamma \text{\$}\vdash ((\varphi \# \Phi) @ \Psi)$   
**from this obtain**  $\Sigma$  **where**  
 $\Sigma: \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma$   
 $\text{map} (\text{uncurry } (\sqcup)) \Sigma \text{\$}\vdash \varphi$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus (\text{map snd } \Sigma) \text{\$}\vdash (\Phi @ \Psi)$   
**(is ? $\Gamma_0 \text{\$}\vdash (\Phi @ \Psi)$ )**

**by auto**

**from this(3) obtain**  $\Delta$  **where**  
 $\Delta: \text{mset} (\text{map snd } \Delta) \subseteq\# \text{mset } ?\Gamma_0$   
 $\text{map} (\text{uncurry } (\sqcup)) \Delta \text{\$}\vdash \Phi$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Delta @ ?\Gamma_0 \ominus (\text{map snd } \Delta) \text{\$}\vdash \Psi$

**using Cons**

by *auto*  
 let  $?\Sigma' = \mathfrak{J} \Sigma \Delta$   
 have  $\text{map } (\text{uncurry } (\sqcup)) \text{ } ?\Sigma' \text{ } \$\vdash (\varphi \# \Phi)$   
   using  $\Delta(1) \Delta(2) \Sigma(2)$  *merge-witness-cons-measure-deduction* by *blast*  
 moreover have  $\text{mset } (\text{map } \text{snd } ?\Sigma') \subseteq\# \text{mset } \Gamma$   
   using  $\Delta(1) \Sigma(1)$  *merge-witness-msub-intro* by *blast*  
 moreover have  $\text{map } (\text{uncurry } (\rightarrow)) \text{ } ?\Sigma' \text{ } @ \Gamma \ominus \text{map } \text{snd } ?\Sigma' \text{ } \$\vdash \Psi$   
   using  $\Delta(1) \Delta(3)$  *merge-witness-measure-deduction-intro* by *blast*  
 ultimately show  
    $\exists \Sigma. \text{mset } (\text{map } \text{snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
      $\text{map } (\text{uncurry } (\sqcup)) \Sigma \text{ } \$\vdash (\varphi \# \Phi) \wedge$   
      $\text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map } \text{snd } \Sigma \text{ } \$\vdash \Psi$   
 by *fast*  
 next  
 assume  $\exists \Sigma. \text{mset } (\text{map } \text{snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
    $\text{map } (\text{uncurry } (\sqcup)) \Sigma \text{ } \$\vdash (\varphi \# \Phi) \wedge$   
    $\text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map } \text{snd } \Sigma \text{ } \$\vdash \Psi$   
 from *this* obtain  $\Delta$  where  $\Delta$ :  
    $\text{mset } (\text{map } \text{snd } \Delta) \subseteq\# \text{mset } \Gamma$   
    $\text{map } (\text{uncurry } (\sqcup)) \Delta \text{ } \$\vdash (\varphi \# \Phi)$   
    $\text{map } (\text{uncurry } (\rightarrow)) \Delta @ \Gamma \ominus \text{map } \text{snd } \Delta \text{ } \$\vdash \Psi$   
 by *auto*  
 from *this* obtain  $\Sigma$  where  $\Sigma$ :  
    $\text{mset } (\text{map } \text{snd } \Sigma) \subseteq\# \text{mset } (\text{map } (\text{uncurry } (\sqcup)) \Delta)$   
    $\text{map } (\text{uncurry } (\sqcup)) \Sigma \text{ } \$\vdash \varphi$   
    $\text{map } (\text{uncurry } (\rightarrow)) \Sigma @ (\text{map } (\text{uncurry } (\sqcup)) \Delta) \ominus \text{map } \text{snd } \Sigma \text{ } \$\vdash \Phi$   
 by *auto*  
 let  $?\Omega = \mathfrak{P} \Sigma \Delta$   
 let  $?\Xi = \mathfrak{Q} \Sigma \Delta$   
 let  $?T_0 = \text{map } (\text{uncurry } (\rightarrow)) \text{ } ?\Omega @ \Gamma \ominus \text{map } \text{snd } ?\Omega$   
 let  $?T_1 = \text{map } (\text{uncurry } (\rightarrow)) \text{ } ?\Xi @ ?T_0 \ominus \text{map } \text{snd } ?\Xi$   
 have  $\text{mset } (\Gamma \ominus \text{map } \text{snd } \Delta) = \text{mset } (?T_0 \ominus \text{map } \text{snd } ?\Xi)$   
   using  $\Delta(1) \Sigma(1)$  *recoverWitnesses-mset-equiv* by *blast*  
 hence  $(\Gamma \ominus \text{map } \text{snd } \Delta) \preceq (?T_0 \ominus \text{map } \text{snd } ?\Xi)$   
   by (*simp add: msub-stronger-theory-intro*)  
 hence  $?T_1 \text{ } \$\vdash \Psi$   
   using  $\Delta(3)$  *measure-stronger-theory-left-monotonic*  
     *stronger-theory-combine*  
     *recover-witness-B-right-stronger-theory*  
 by *blast*  
 moreover  
 have  $\text{mset } (\text{map } \text{snd } ?\Xi) \subseteq\# \text{mset } ?T_0$   
   using  $\Sigma(1) \Delta(1)$  *recover-witness-B-mset-equiv*  
   by (*simp*,  
     *metis list-subtract-monotonic*  
     *list-subtract-mset-homomorphism*  
     *mset-map*)  
 moreover  
 have  $\text{map } (\text{uncurry } (\sqcup)) \text{ } ?\Xi \text{ } \$\vdash \Phi$

```

    using  $\Sigma(1)$  recover-witness-B-stronger-theory
       $\Sigma(3)$  measure-stronger-theory-left-monotonic by blast
  ultimately have  $?T_0 \ \$\vdash (\Phi @ \Psi)$ 
    using Cons by fast
  moreover
  have  $mset (map\ snd\ ?\Omega) \subseteq\# mset (map\ snd\ \Delta)$ 
    using  $\Sigma(1)$  recover-witness-A-mset-equiv
    by (simp, metis mset-subset-eq-add-left)
  hence  $mset (map\ snd\ ?\Omega) \subseteq\# mset\ \Gamma$  using  $\Delta(1)$  by simp
  moreover
  have  $map (uncurry (\sqcup))\ ?\Omega \vdash \varphi$ 
    using  $\Sigma(2)$ 
      recover-witness-A-left-stronger-theory
      stronger-theory-deduction-monotonic
    by blast
  ultimately show  $\Gamma \ \$\vdash ((\varphi \# \Phi) @ \Psi)$ 
    by (simp, blast)
  qed
}
then show ?case by metis
  qed
thus ?thesis by blast
  qed

```

**lemma** (in *classical-logic*) *measure-list-deduction-antitonic*:

```

  assumes  $\Gamma \ \$\vdash \Psi$ 
    and  $\Psi \vdash \varphi$ 
  shows  $\Gamma \vdash \varphi$ 
  using assms
proof (induct  $\Psi$  arbitrary:  $\Gamma \ \varphi$ )
  case Nil
  then show ?case
    using list-deduction-weaken
    by simp
next
  case (Cons  $\psi \ \Psi$ )
  hence  $\Psi \vdash \psi \rightarrow \varphi$ 
    using list-deduction-theorem by blast
  from  $\langle \Gamma \ \$\vdash (\psi \# \Psi) \rangle$  obtain  $\Sigma$  where  $\Sigma$ :
     $mset (map\ snd\ \Sigma) \subseteq\# mset\ \Gamma$ 
     $map (uncurry (\sqcup))\ \Sigma \vdash \psi$ 
     $map (uncurry (\rightarrow))\ \Sigma @ \Gamma \ominus map\ snd\ \Sigma \ \$\vdash \Psi$ 
  by auto
  hence  $\Gamma \vdash \psi \rightarrow \varphi$ 
  using
    measure-stronger-theory-left-monotonic
    witness-stronger-theory
     $\langle \Psi \vdash \psi \rightarrow \varphi \rangle$ 
    Cons

```

```

  by blast
  moreover
  have  $\Gamma \vdash \psi$ 
  using  $\Sigma(1) \Sigma(2)$ 
    stronger-theory-deduction-monotonic
    witness-weaker-theory
  by blast
  ultimately show ?case using list-deduction-modus-ponens by auto
qed

```

Finally, we may establish that  $(\mathbb{S}\vdash)$  is transitive.

```

theorem (in classical-logic) measure-transitive:
  assumes  $\Gamma \mathbb{S}\vdash \Lambda$ 
    and  $\Lambda \mathbb{S}\vdash \Delta$ 
  shows  $\Gamma \mathbb{S}\vdash \Delta$ 
  using assms
proof (induct  $\Delta$  arbitrary:  $\Gamma \Lambda$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\delta \Delta$ )
  from this obtain  $\Sigma$  where  $\Sigma$ :
    mset (map snd  $\Sigma$ )  $\subseteq\#$  mset  $\Lambda$ 
    map (uncurry ( $\sqcup$ ))  $\Sigma \vdash \delta$ 
    map (uncurry ( $\rightarrow$ ))  $\Sigma @ \Lambda \ominus$  map snd  $\Sigma \mathbb{S}\vdash \Delta$ 
  by auto
  hence  $\Gamma \mathbb{S}\vdash$  (map (uncurry ( $\sqcup$ ))  $\Sigma @$  map (uncurry ( $\rightarrow$ ))  $\Sigma @ \Lambda \ominus$  (map snd
 $\Sigma$ ))
  using Cons measure-witness-right-split
  by simp
  from this obtain  $\Psi$  where  $\Psi$ :
    mset (map snd  $\Psi$ )  $\subseteq\#$  mset  $\Gamma$ 
    map (uncurry ( $\sqcup$ ))  $\Psi \mathbb{S}\vdash$  map (uncurry ( $\sqcup$ ))  $\Sigma$ 
    map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus$  map snd  $\Psi \mathbb{S}\vdash$  (map (uncurry ( $\rightarrow$ ))  $\Sigma @ \Lambda \ominus$ 
map snd  $\Sigma$ )
  using measure-deduction-generalized-witness
  by fastforce
  have map (uncurry ( $\rightarrow$ ))  $\Psi @ \Gamma \ominus$  map snd  $\Psi \mathbb{S}\vdash \Delta$ 
  using  $\Sigma(3) \Psi(3) Cons$ 
  by auto
  moreover
  have map (uncurry ( $\sqcup$ ))  $\Psi \vdash \delta$ 
  using  $\Psi(2) \Sigma(2)$  measure-list-deduction-antitonic
  by blast
  ultimately show ?case
  using  $\Psi(1)$ 
  by fastforce
qed

```

## 2.6 Measure Deduction Cancellation Rules

In this chapter we go over how to cancel formulae occurring in measure deduction judgements.

The first observation is that tautologies can always be canceled on either side of the turnstile.

**lemma** (in *classical-logic*) *measure-tautology-right-cancel*:

```

assumes  $\vdash \varphi$ 
shows  $\Gamma \ \$\vdash (\varphi \# \Phi) = \Gamma \ \$\vdash \Phi$ 
proof (rule iffI)
assume  $\Gamma \ \$\vdash (\varphi \# \Phi)$ 
from this obtain  $\Sigma$  where  $\Sigma$ :
   $mset (map \ snd \ \Sigma) \subseteq\# \ mset \ \Gamma$ 
   $map (uncurry (\sqcup)) \ \Sigma \ :\vdash \ \varphi$ 
   $map (uncurry (\rightarrow)) \ \Sigma \ @ \ \Gamma \ \ominus \ map \ snd \ \Sigma \ \$\vdash \ \Phi$ 
by auto
thus  $\Gamma \ \$\vdash \Phi$ 
using measure-stronger-theory-left-monotonic
       witness-stronger-theory
by blast
next
assume  $\Gamma \ \$\vdash \Phi$ 
hence  $map (uncurry (\rightarrow)) \ [] \ @ \ \Gamma \ \ominus \ map \ snd \ [] \ \$\vdash \Phi$ 
        $mset (map \ snd \ []) \subseteq\# \ mset \ \Gamma$ 
        $map (uncurry (\sqcup)) \ [] \ :\vdash \ \varphi$ 
using assms
by simp+
thus  $\Gamma \ \$\vdash (\varphi \# \Phi)$ 
using measure-deduction.simps(2)
by blast
qed

```

**lemma** (in *classical-logic*) *measure-tautology-left-cancel* [*simp*]:

```

assumes  $\vdash \gamma$ 
shows  $(\gamma \# \Gamma) \ \$\vdash \Phi = \Gamma \ \$\vdash \Phi$ 
proof (rule iffI)
assume  $(\gamma \# \Gamma) \ \$\vdash \Phi$ 
moreover have  $\Gamma \ \$\vdash \Gamma$ 
  by (simp add: stronger-theory-to-measure-deduction)
hence  $\Gamma \ \$\vdash (\gamma \# \Gamma)$ 
using assms measure-tautology-right-cancel
by simp
ultimately show  $\Gamma \ \$\vdash \Phi$ 
using measure-transitive by blast
next
assume  $\Gamma \ \$\vdash \Phi$ 
moreover have  $mset \ \Gamma \subseteq\# \ mset (\gamma \# \Gamma)$ 
by simp

```



```

hence ( $\gamma \# \Gamma$ )  $\$ \vdash \Gamma$ 
  using msub-stronger-theory-intro
    stronger-theory-to-measure-deduction
  by blast
ultimately show ( $\gamma \# \Gamma$ )  $\$ \vdash \Phi$ 
  using measure-transitive by blast
qed

```

**lemma** (in *classical-logic*) *measure-deduction-one-collapse*:

```

 $\Gamma \ \$ \vdash [\varphi] = \Gamma \ : \vdash \varphi$ 
proof (rule iffI)
  assume  $\Gamma \ \$ \vdash [\varphi]$ 
  from this obtain  $\Sigma$  where
     $\Sigma$ :  $mset (map \text{snd } \Sigma) \subseteq \# mset \Gamma$ 
     $map (uncurry (\sqcup)) \Sigma \ : \vdash \varphi$ 
  by auto
  hence  $map (uncurry (\sqcup)) \Sigma \preceq \Gamma$ 
  using witness-weaker-theory by blast
  thus  $\Gamma \ : \vdash \varphi$  using  $\Sigma(2)$ 
  using stronger-theory-deduction-monotonic by blast
next
  assume  $\Gamma \ : \vdash \varphi$ 
  let  $?\Sigma = map (\lambda \gamma. (\perp, \gamma)) \Gamma$ 
  have  $\Gamma \preceq map (uncurry (\sqcup)) ?\Sigma$ 
  proof (induct  $\Gamma$ )
    case Nil
    then show ?case by simp
  next
  case (Cons  $\gamma \Gamma$ )
  have  $\vdash (\perp \sqcup \gamma) \rightarrow \gamma$ 
  unfolding disjunction-def
  using ex-falso-quodlibet modus-ponens excluded-middle-elimination
  by blast
  then show ?case using Cons
  by (simp add: stronger-theory-left-right-cons)
qed
  hence  $map (uncurry (\sqcup)) ?\Sigma \ : \vdash \varphi$ 
  using  $\langle \Gamma \ : \vdash \varphi \rangle$  stronger-theory-deduction-monotonic by blast
  moreover have  $mset (map \text{snd } ?\Sigma) \subseteq \# mset \Gamma$  by (induct  $\Gamma$ , simp+)
  ultimately show  $\Gamma \ \$ \vdash [\varphi]$ 
  using measure-deduction.simps(1)
    measure-deduction.simps(2)
  by blast
qed

```

*Split cases*, which are occurrences of  $\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \dots$ , also cancel and simplify to just  $\varphi \# \dots$ . We previously established  $\Gamma \ \$ \vdash \psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Phi = \Gamma \ \$ \vdash \varphi \# \Phi$  as part of the proof of transitivity.

**lemma** (in *classical-logic*) *measure-formula-left-split*:  
 $\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma \ \$\vdash \Phi = \varphi \# \Gamma \ \$\vdash \Phi$   
**proof** (*rule iffI*)  
**assume**  $\varphi \# \Gamma \ \$\vdash \Phi$   
**have**  $\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma \ \$\vdash (\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma)$   
**using** *stronger-theory-to-measure-deduction*  
*stronger-theory-reflexive*  
**by** *blast*  
**hence**  $\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma \ \$\vdash (\varphi \# \Gamma)$   
**using** *measure-formula-right-split* **by** *blast*  
**with**  $\langle \varphi \# \Gamma \ \$\vdash \Phi \rangle$  **show**  $\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma \ \$\vdash \Phi$   
**using** *measure-transitive* **by** *blast*  
**next**  
**assume**  $\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma \ \$\vdash \Phi$   
**have**  $\varphi \# \Gamma \ \$\vdash (\varphi \# \Gamma)$   
**using** *stronger-theory-to-measure-deduction*  
*stronger-theory-reflexive*  
**by** *blast*  
**hence**  $\varphi \# \Gamma \ \$\vdash (\psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma)$   
**using** *measure-formula-right-split* **by** *blast*  
**with**  $\langle \psi \sqcup \varphi \# \psi \rightarrow \varphi \# \Gamma \ \$\vdash \Phi \rangle$  **show**  $\varphi \# \Gamma \ \$\vdash \Phi$   
**using** *measure-transitive* **by** *blast*  
**qed**

**lemma** (in *classical-logic*) *measure-witness-left-split* [*simp*]:  
**assumes**  $mset (map \text{snd } \Sigma) \subseteq\# mset \Gamma$   
**shows**  $(map (uncurry (\sqcup)) \Sigma @ map (uncurry (\rightarrow)) \Sigma @ \Gamma \ominus (map \text{snd } \Sigma)) \ \$\vdash \Phi = \Gamma \ \$\vdash \Phi$   
**using** *assms*  
**proof** (*induct*  $\Sigma$  *arbitrary*:  $\Gamma$ )  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons*  $\sigma \Sigma$ )  
**let**  $?x = \text{fst } \sigma$   
**let**  $?y = \text{snd } \sigma$   
**let**  $?T_0 = map (uncurry (\sqcup)) \Sigma @ map (uncurry (\rightarrow)) \Sigma @ \Gamma \ominus map \text{snd } (\sigma \# \Sigma)$   
**let**  $?T' = map (uncurry (\sqcup)) (\sigma \# \Sigma) @ map (uncurry (\rightarrow)) (\sigma \# \Sigma) @ \Gamma \ominus map \text{snd } (\sigma \# \Sigma)$   
**assume**  $mset (map \text{snd } (\sigma \# \Sigma)) \subseteq\# mset \Gamma$   
**hence** *A*:  $add\text{-}mset (\text{snd } \sigma) (image\text{-}mset \text{snd } (mset \Sigma)) \subseteq\# mset \Gamma$  **by** *simp*  
**hence** *B*:  $image\text{-}mset \text{snd } (mset \Sigma) + (mset \Gamma - image\text{-}mset \text{snd } (mset \Sigma))$   
 $= add\text{-}mset (\text{snd } \sigma) (image\text{-}mset \text{snd } (mset \Sigma))$   
 $+ (mset \Gamma - add\text{-}mset (\text{snd } \sigma) (image\text{-}mset \text{snd } (mset \Sigma)))$   
**by** (*metis* (*no-types*) *mset-subset-eq-insertD* *subset-mset.add-diff-inverse* *subset-mset-def*)  
**have**  $\{\#x \rightarrow y. (x, y) \in\# mset \Sigma\# \}$   
 $+ mset \Gamma - add\text{-}mset (\text{snd } \sigma) (image\text{-}mset \text{snd } (mset \Sigma))$

$= \{ \#x \rightarrow y. (x, y) \in \# \text{mset } \Sigma \# \}$   
 $+ (\text{mset } \Gamma - \text{add-mset } (\text{snd } \sigma) (\text{image-mset } \text{snd } (\text{mset } \Sigma)))$   
**using** *A subset-mset.diff-add-assoc* **by** *blast*  
**hence**  $\{ \#x \rightarrow y. (x, y) \in \# \text{mset } \Sigma \# \} + (\text{mset } \Gamma - \text{image-mset } \text{snd } (\text{mset } \Sigma))$   
 $= \text{add-mset } (\text{snd } \sigma) (\{ \#x \rightarrow y. (x, y) \in \# \text{mset } \Sigma \# \})$   
 $+ \text{mset } \Gamma - \text{add-mset } (\text{snd } \sigma) (\text{image-mset } \text{snd } (\text{mset } \Sigma))$   
**using** *B* **by** *auto*  
**hence** *C*:  
 $\text{mset } (\text{map } \text{snd } \Sigma) \subseteq \# \text{mset } \Gamma$   
 $\text{mset } (\text{map } (\text{uncurry } (\sqcup)) \Sigma @ \text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map } \text{snd } \Sigma)$   
 $= \text{mset } (? \gamma \# ? \Gamma_0)$   
**using**  $\langle \text{mset } (\text{map } \text{snd } (\sigma \# \Sigma)) \subseteq \# \text{mset } \Gamma \rangle$   
*subset-mset.dual-order.trans*  
**by** *(fastforce+)*  
**hence**  $\Gamma \text{ \$} \vdash \Phi = (? \chi \sqcup ? \gamma \# ? \chi \rightarrow ? \gamma \# ? \Gamma_0) \text{ \$} \vdash \Phi$   
**proof** –  
**have**  $\forall \Gamma \Delta. \neg \text{mset } (\text{map } \text{snd } \Sigma) \subseteq \# \text{mset } \Gamma$   
 $\vee \neg \Gamma \text{ \$} \vdash \Phi$   
 $\vee \neg \text{mset } (\text{map } (\text{uncurry } (\sqcup)) \Sigma$   
 $\quad @ \text{map } (\text{uncurry } (\rightarrow)) \Sigma$   
 $\quad @ \Gamma \ominus \text{map } \text{snd } \Sigma)$   
 $\subseteq \# \text{mset } \Delta$   
 $\vee \Delta \text{ \$} \vdash \Phi$   
**using** *Cons.hyps measure-msub-left-monotonic* **by** *blast*  
**moreover**  
{  
**assume**  $\neg \Gamma \text{ \$} \vdash \Phi$   
**then have**  $\exists \Delta. \text{mset } (\text{snd } \sigma \# \text{map } (\text{uncurry } (\sqcup)) \Sigma$   
 $\quad @ \text{map } (\text{uncurry } (\rightarrow)) \Sigma$   
 $\quad @ \Gamma \ominus \text{map } \text{snd } (\sigma \# \Sigma))$   
 $\subseteq \# \text{mset } \Delta$   
 $\wedge \neg \Gamma \text{ \$} \vdash \Phi$   
 $\wedge \neg \Delta \text{ \$} \vdash \Phi$   
**by** *(metis (no-types) Cons.hyps C subset-mset.dual-order.refl)*  
**then have** *?thesis*  
**using** *measure-formula-left-split measure-msub-left-monotonic* **by** *blast*  
}  
**ultimately show** *?thesis*  
**by** *(metis (full-types) C measure-formula-left-split subset-mset.dual-order.refl)*  
**qed**  
**moreover**  
**have**  $(\text{uncurry } (\sqcup)) = (\lambda \psi. \text{fst } \psi \sqcup \text{snd } \psi)$   
 $(\text{uncurry } (\rightarrow)) = (\lambda \psi. \text{fst } \psi \rightarrow \text{snd } \psi)$   
**by** *fastforce+*  
**hence**  $\text{mset } ? \Gamma' = \text{mset } (? \chi \sqcup ? \gamma \# ? \chi \rightarrow ? \gamma \# ? \Gamma_0)$   
**by** *fastforce*  
**hence**  $(? \chi \sqcup ? \gamma \# ? \chi \rightarrow ? \gamma \# ? \Gamma_0) \text{ \$} \vdash \Phi = ? \Gamma' \text{ \$} \vdash \Phi$   
**by** *(metis*  
*(mono-tags, lifting)*)

```

      measure-msub-left-monotonic
      subset-mset.dual-order.refl)
ultimately have  $\Gamma \text{\$}\vdash \Phi = ?\Gamma' \text{\$}\vdash \Phi$ 
  by fastforce
then show ?case by blast
qed

```

We now have enough to establish the cancellation rule for ( $\text{\$}\vdash$ ).

**lemma** (in *classical-logic*) *measure-cancel*:  $(\Delta @ \Gamma) \text{\$}\vdash (\Delta @ \Phi) = \Gamma \text{\$}\vdash \Phi$

**proof** –

```

{
  fix  $\Delta \Gamma \Phi$ 
  assume  $\Gamma \text{\$}\vdash \Phi$ 
  hence  $(\Delta @ \Gamma) \text{\$}\vdash (\Delta @ \Phi)$ 
  proof (induct  $\Delta$ )
    case Nil
      then show ?case by simp
    next
      case (Cons  $\delta \Delta$ )
        let  $?\Sigma = [(\delta, \delta)]$ 
        have map (uncurry ( $\sqcup$ ))  $?\Sigma \text{\$}\vdash \delta$ 
          unfolding disjunction-def list-duction-def
          by (simp add: Peirces-law)
        moreover have mset (map snd  $?\Sigma$ )  $\subseteq\#$  mset ( $\delta \# \Delta$ ) by simp
        moreover have map (uncurry ( $\rightarrow$ ))  $?\Sigma @ ((\delta \# \Delta) @ \Gamma) \ominus$  map snd  $?\Sigma \text{\$}\vdash$ 
          ( $\Delta @ \Phi$ )
          using Cons
          by (simp add: trivial-implication)
        moreover have map snd  $[(\delta, \delta)] = [\delta]$  by force
        ultimately show ?case
          by (metis (no-types) measure-duction.simps(2)
              append-Cons
              list.set-intros(1)
              mset.simps(1)
              mset.simps(2)
              mset-subset-eq-single
              set-mset-mset)
  qed
} note forward-direction = this
{
  assume  $(\Delta @ \Gamma) \text{\$}\vdash (\Delta @ \Phi)$ 
  hence  $\Gamma \text{\$}\vdash \Phi$ 
  proof (induct  $\Delta$ )
    case Nil
      then show ?case by simp
    next
      case (Cons  $\delta \Delta$ )
        have mset  $((\delta \# \Delta) @ \Phi) =$  mset  $((\Delta @ \Phi) @ [\delta])$  by simp
        with Cons.premis have  $((\delta \# \Delta) @ \Gamma) \text{\$}\vdash ((\Delta @ \Phi) @ [\delta])$ 

```

```

    by (metis measure-msub-weaken
        subset-mset.dual-order.refl)
from this obtain  $\Sigma$  where  $\Sigma$ :
  mset (map snd  $\Sigma$ )  $\subseteq\#$  mset (( $\delta \# \Delta$ ) @  $\Gamma$ )
  map (uncurry ( $\sqcup$ ))  $\Sigma$   $\$ \vdash$  ( $\Delta$  @  $\Phi$ )
  map (uncurry ( $\rightarrow$ ))  $\Sigma$  @ (( $\delta \# \Delta$ ) @  $\Gamma$ )  $\ominus$  map snd  $\Sigma$   $\$ \vdash$  [ $\delta$ ]
  by (metis append-assoc measure-deduction-generalized-witness)
show ?case
proof (cases find ( $\lambda \sigma. \text{snd } \sigma = \delta$ )  $\Sigma = \text{None}$ )
  case True
  hence  $\delta \notin \text{set } (\text{map snd } \Sigma)$ 
  proof (induct  $\Sigma$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\sigma$   $\Sigma$ )
    then show ?case by (cases snd  $\sigma = \delta$ , simp+)
  qed
with  $\Sigma(1)$  have mset (map snd  $\Sigma$ )  $\subseteq\#$  mset ( $\Delta$  @  $\Gamma$ )
  by (simp, metis add-mset-add-single
      diff-single-trivial
      mset-map
      set-mset-mset
      subset-eq-diff-conv)
thus ?thesis
  using measure-stronger-theory-left-monotonic
      witness-weaker-theory
      Cons.hyps  $\Sigma(2)$ 
  by blast
next
case False
from this obtain  $\sigma$   $\chi$  where
   $\sigma: \sigma = (\chi, \delta)$ 
   $\sigma \in \text{set } \Sigma$ 
  using find-Some-predicate
      find-Some-set-membership
  by fastforce
let  $?\Sigma' = \text{remove1 } \sigma \Sigma$ 
let  $?\Sigma_A = \text{map } (\text{uncurry } (\sqcup)) \ ?\Sigma'$ 
let  $?\Sigma_B = \text{map } (\text{uncurry } (\rightarrow)) \ ?\Sigma'$ 
have mset  $\Sigma = \text{mset } (?\Sigma' @ [(\chi, \delta)])$ 
  mset  $\Sigma = \text{mset } ((\chi, \delta) \# ?\Sigma')$ 
  using  $\sigma$  by simp+
hence mset (map (uncurry ( $\sqcup$ ))  $\Sigma$ ) = mset (map (uncurry ( $\sqcup$ )) ( ?\Sigma' @ [(\chi,
 $\delta$ )]))
  mset (map snd  $\Sigma$ ) = mset (map snd (( $\chi, \delta$ ) # ?\Sigma'))
  mset (map (uncurry ( $\rightarrow$ ))  $\Sigma$ ) = mset (map (uncurry ( $\rightarrow$ )) (( $\chi, \delta$ ) #
? $\Sigma'$ ))
  by (metis mset-map)+

```

**hence**  $mset (map (uncurry (\sqcup)) \Sigma) = mset (? \Sigma_A @ [\chi \sqcup \delta])$   
 $mset (map (uncurry (\rightarrow)) \Sigma @ ((\delta \# \Delta) @ \Gamma) \ominus map\ snd\ \Sigma)$   
 $= mset (\chi \rightarrow \delta \# ? \Sigma_B @ (\Delta @ \Gamma) \ominus map\ snd\ ? \Sigma')$   
**by** *simp+*  
**hence**  
 $? \Sigma_A @ [\chi \sqcup \delta] \$\vdash (\Delta @ \Phi)$   
 $\chi \rightarrow \delta \# (? \Sigma_B @ (\Delta @ \Gamma) \ominus map\ snd\ ? \Sigma') \$\vdash [\delta]$   
**using**  $\Sigma(2)\ \Sigma(3)$   
**by** (*metis measure-msub-left-monotonic subset-mset.dual-order.refl, simp*)  
**moreover**  
**have**  $\vdash ((\chi \rightarrow \delta) \rightarrow \delta) \rightarrow (\chi \sqcup \delta)$   
**unfolding** *disjunction-def*  
**using** *modus-ponens*  
*pseudo-scotus*  
*flip-hypothetical-syllogism*  
**by** *blast*  
**ultimately have**  $(? \Sigma_A @ ? \Sigma_B @ (\Delta @ \Gamma) \ominus map\ snd\ ? \Sigma') \$\vdash (\Delta @ \Phi)$   
**using** *measure-deduction-one-collapse*  
*list-deduction-theorem*  
*list-deduction-modus-ponens*  
*list-deduction-weaken*  
*forward-direction*  
*measure-transitive*  
**by** *meson*  
**moreover**  
**have**  $\delta = snd\ \sigma$   
 $snd\ \sigma \in set (map\ snd\ \Sigma)$   
**by** (*simp add: \sigma(1), simp add: \sigma(2)*)  
**with**  $\Sigma(1)$  **have**  $mset (map\ snd (remove1\ \sigma\ \Sigma)) \subseteq\# mset (remove1\ \delta ((\delta$   
 $\# \Delta) @ \Gamma))$   
**by** (*metis insert-DiffM*  
*insert-subset-eq-iff*  
*mset-remove1*  
 $\sigma(1)\ \sigma(2)$   
*remove1-pairs-list-projections-snd*  
*set-mset-mset*)  
**hence**  $mset (map\ snd (remove1\ \sigma\ \Sigma)) \subseteq\# mset (\Delta @ \Gamma)$  **by** *simp*  
**ultimately show** *?thesis*  
**using** *measure-witness-left-split Cons.hyps*  
**by** *blast*  
**qed**  
**qed**  
**}**  
**with** *forward-direction* **show** *?thesis* **by** *auto*  
**qed**

**lemma** (in *classical-logic*) *measure-biconditional-cancel*:

**assumes**  $\vdash \gamma \leftrightarrow \varphi$

**shows**  $(\gamma \# \Gamma) \$\vdash (\varphi \# \Phi) = \Gamma \$\vdash \Phi$

**proof** –  
**from** *assms* **have**  $(\gamma \# \Phi) \preceq (\varphi \# \Phi) (\varphi \# \Phi) \preceq (\gamma \# \Phi)$   
**unfolding** *biconditional-def*  
**by** (*simp add: stronger-theory-left-right-cons*)  
**hence**  $(\gamma \# \Phi) \$\vdash (\varphi \# \Phi)$   
 $(\varphi \# \Phi) \$\vdash (\gamma \# \Phi)$   
**using** *stronger-theory-to-measure-deduction* **by** *blast+*  
**moreover**  
**have**  $\Gamma \$\vdash \Phi = (\gamma \# \Gamma) \$\vdash (\gamma \# \Phi)$   
**by** (*metis append-Cons append-Nil measure-cancel*)  
**ultimately**  
**have**  $\Gamma \$\vdash \Phi \implies \gamma \# \Gamma \$\vdash (\varphi \# \Phi)$   
 $\gamma \# \Gamma \$\vdash (\varphi \# \Phi) \implies \Gamma \$\vdash \Phi$   
**using** *measure-transitive* **by** *blast+*  
**thus** *?thesis* **by** *blast*  
**qed**

## 2.7 Measure Deduction Substitution Rules

Just like conventional deduction, if two formulae are equivalent then they may be substituted for one another.

**lemma** (in *classical-logic*) *right-measure-sub*:

**assumes**  $\vdash \varphi \leftrightarrow \psi$   
**shows**  $\Gamma \$\vdash (\varphi \# \Phi) = \Gamma \$\vdash (\psi \# \Phi)$   
**proof** –  
**have**  $\Gamma \$\vdash (\varphi \# \Phi) = (\psi \# \Gamma) \$\vdash (\psi \# \varphi \# \Phi)$   
**using** *measure-cancel* [**where**  $\Delta=[\psi]$  **and**  $\Gamma=\Gamma$  **and**  $\Phi=\varphi \# \Phi$ ] **by** *simp*  
**also have**  $\dots = (\psi \# \Gamma) \$\vdash (\varphi \# \psi \# \Phi)$   
**using** *measure-cons-cons-right-permute* **by** *blast*  
**also have**  $\dots = \Gamma \$\vdash (\psi \# \Phi)$   
**using** *assms biconditional-symmetry-rule measure-biconditional-cancel* **by** *blast*  
**finally show** *?thesis* .  
**qed**

**lemma** (in *classical-logic*) *left-measure-sub*:

**assumes**  $\vdash \gamma \leftrightarrow \chi$   
**shows**  $(\gamma \# \Gamma) \$\vdash \Phi = (\chi \# \Gamma) \$\vdash \Phi$   
**proof** –  
**have**  $(\gamma \# \Gamma) \$\vdash \Phi = (\chi \# \gamma \# \Gamma) \$\vdash (\chi \# \Phi)$   
**using** *measure-cancel* [**where**  $\Delta=[\chi]$  **and**  $\Gamma=(\gamma \# \Gamma)$  **and**  $\Phi=\Phi$ ] **by** *simp*  
**also have**  $\dots = (\gamma \# \chi \# \Gamma) \$\vdash (\chi \# \Phi)$   
**using**  
*measure-cons-cons-right-permute*  
*stronger-theory-to-measure-deduction*  
*measure-transitive*  
*stronger-theory-reflexive*  
**by** *blast*  
**also have**  $\dots = (\chi \# \Gamma) \$\vdash \Phi$

using *assms biconditional-symmetry-rule measure-biconditional-cancel* **by blast**  
 finally show *?thesis* .  
 qed

## 2.8 Measure Deduction Sum Rules

We next establish analogues of the rule in probability that  $\mathcal{P} \alpha + \mathcal{P} \beta = \mathcal{P} (\alpha \sqcup \beta) + \mathcal{P} (\alpha \sqcap \beta)$ . This equivalence holds for both sides of the  $(\$ \vdash)$  turnstile.

**lemma** (in *classical-logic*) *right-measure-sum-rule*:

$$\Gamma \$ \vdash (\alpha \# \beta \# \Phi) = \Gamma \$ \vdash (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \Phi)$$

**proof** –

$$\text{have } A: \text{mset } (\alpha \sqcup \beta \# \beta \rightarrow \alpha \# \beta \# \Phi) = \text{mset } (\beta \rightarrow \alpha \# \beta \# \alpha \sqcup \beta \# \Phi)$$

by *simp*

$$\text{have } B: \vdash (\beta \rightarrow \alpha) \leftrightarrow (\beta \rightarrow (\alpha \sqcap \beta))$$

**proof** –

$$\text{let } ?\varphi = (\langle \beta \rangle \rightarrow \langle \alpha \rangle) \leftrightarrow (\langle \beta \rangle \rightarrow (\langle \alpha \rangle \sqcap \langle \beta \rangle))$$

$$\text{have } \forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi \text{ by } \textit{fastforce}$$

$$\text{hence } \vdash (\langle ?\varphi \rangle) \text{ using } \textit{propositional-semantic} \text{ by } \textit{blast}$$

$$\text{thus } ?\textit{thesis} \text{ by } \textit{simp}$$

qed

$$\text{have } C: \vdash \beta \leftrightarrow (\beta \sqcup (\alpha \sqcap \beta))$$

**proof** –

$$\text{let } ?\varphi = \langle \beta \rangle \leftrightarrow (\langle \beta \rangle \sqcup (\langle \alpha \rangle \sqcap \langle \beta \rangle))$$

$$\text{have } \forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi \text{ by } \textit{fastforce}$$

$$\text{hence } \vdash (\langle ?\varphi \rangle) \text{ using } \textit{propositional-semantic} \text{ by } \textit{blast}$$

$$\text{thus } ?\textit{thesis} \text{ by } \textit{simp}$$

qed

$$\text{have } \Gamma \$ \vdash (\alpha \# \beta \# \Phi) = \Gamma \$ \vdash (\beta \sqcup \alpha \# \beta \rightarrow \alpha \# \beta \# \Phi)$$

$$\text{using } \textit{measure-formula-right-split} \text{ by } \textit{blast}$$

$$\text{also have } \dots = \Gamma \$ \vdash (\alpha \sqcup \beta \# \beta \rightarrow \alpha \# \beta \# \Phi)$$

$$\text{using } \textit{disjunction-commutativity right-measure-sub} \text{ by } \textit{blast}$$

$$\text{also have } \dots = \Gamma \$ \vdash (\beta \rightarrow \alpha \# \beta \# \alpha \sqcup \beta \# \Phi)$$

$$\text{by } (\textit{metis } A \textit{ measure-msub-weaken subset-mset.dual-order.refl})$$

$$\text{also have } \dots = \Gamma \$ \vdash (\beta \rightarrow (\alpha \sqcap \beta) \# \beta \# \alpha \sqcup \beta \# \Phi)$$

$$\text{using } B \textit{ right-measure-sub} \text{ by } \textit{blast}$$

$$\text{also have } \dots = \Gamma \$ \vdash (\beta \# \beta \rightarrow (\alpha \sqcap \beta) \# \alpha \sqcup \beta \# \Phi)$$

$$\text{using } \textit{measure-cons-cons-right-permute} \text{ by } \textit{blast}$$

$$\text{also have } \dots = \Gamma \$ \vdash (\beta \sqcup (\alpha \sqcap \beta) \# \beta \rightarrow (\alpha \sqcap \beta) \# \alpha \sqcup \beta \# \Phi)$$

$$\text{using } C \textit{ right-measure-sub} \text{ by } \textit{blast}$$

$$\text{also have } \dots = \Gamma \$ \vdash (\alpha \sqcap \beta \# \alpha \sqcup \beta \# \Phi)$$

$$\text{using } \textit{measure-formula-right-split} \text{ by } \textit{blast}$$

finally show *?thesis*

$$\text{using } \textit{measure-cons-cons-right-permute} \text{ by } \textit{blast}$$

qed

**lemma** (in *classical-logic*) *left-measure-sum-rule*:

$$(\alpha \# \beta \# \Gamma) \$ \vdash \Phi = (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \Gamma) \$ \vdash \Phi$$



**proof** –

**have**  $\star$ :  $mset (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \alpha \# \beta \# \Gamma) = mset (\alpha \# \beta \# \alpha \sqcup \beta \# \alpha \sqcap \beta \# \Gamma)$  **by** *simp*

**have**  $(\alpha \# \beta \# \Gamma) \$\vdash \Phi = (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \alpha \# \beta \# \Gamma) \$\vdash (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \Phi)$

**using** *measure-cancel* [**where**  $\Delta=[\alpha \sqcup \beta, \alpha \sqcap \beta]$  **and**  $\Gamma=(\alpha \# \beta \# \Gamma)$  **and**  $\Phi=\Phi$ ] **by** *simp*

**also have**  $\dots = (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \alpha \# \beta \# \Gamma) \$\vdash (\alpha \# \beta \# \Phi)$

**using** *right-measure-sum-rule* **by** *blast*

**also have**  $\dots = (\alpha \# \beta \# \alpha \sqcup \beta \# \alpha \sqcap \beta \# \Gamma) \$\vdash (\alpha \# \beta \# \Phi)$

**by** (*metis*  $\star$  *measure-msub-left-monotonic subset-mset.dual-order.refl*)

**also have**  $\dots = (\alpha \sqcup \beta \# \alpha \sqcap \beta \# \Gamma) \$\vdash \Phi$

**using** *measure-cancel* [**where**  $\Delta=[\alpha, \beta]$  **and**  $\Gamma=(\alpha \sqcup \beta \# \alpha \sqcap \beta \# \Gamma)$  **and**  $\Phi=\Phi$ ] **by** *simp*

**finally show** *?thesis* .

**qed**

## 2.9 Measure Deduction Exchange Rule

As we will see, a key result is that we can move formulae from the right hand side of the ( $\$\vdash$ ) turnstile to the left.

We observe a novel logical principle, which we call *exchange*. This principle follows immediately from the split rules and cancellation rules.

**lemma** (in *classical-logic*) *measure-exchange*:

$(\gamma \# \Gamma) \$\vdash (\varphi \# \Phi) = (\varphi \rightarrow \gamma \# \Gamma) \$\vdash (\gamma \rightarrow \varphi \# \Phi)$

**proof** –

**have**  $(\gamma \# \Gamma) \$\vdash (\varphi \# \Phi) = (\varphi \sqcup \gamma \# \varphi \rightarrow \gamma \# \Gamma) \$\vdash (\gamma \sqcup \varphi \# \gamma \rightarrow \varphi \# \Phi)$

**using** *measure-formula-left-split*

*measure-formula-right-split*

**by** *blast+*

**thus** *?thesis*

**using** *measure-biconditional-cancel*

*disjunction-commutativity*

**by** *blast*

**qed**

The exchange rule allows us to prove an analogue of the rule in classical logic that  $\Gamma \vdash \varphi = (\sim \varphi \# \Gamma) \vdash \perp$  for measure deduction.

**theorem** (in *classical-logic*) *measure-negation-swap*:

$\Gamma \$\vdash (\varphi \# \Phi) = (\sim \varphi \# \Gamma) \$\vdash (\perp \# \Phi)$

**proof** –

**have**  $\Gamma \$\vdash (\varphi \# \Phi) = (\perp \# \Gamma) \$\vdash (\perp \# \varphi \# \Phi)$

**by** (*metis* *append-Cons append-Nil measure-cancel*)

**also have**  $\dots = (\perp \# \Gamma) \$\vdash (\varphi \# \perp \# \Phi)$

**using** *measure-cons-cons-right-permute* **by** *blast*

**also have**  $\dots = (\sim \varphi \# \Gamma) \$\vdash (\perp \rightarrow \varphi \# \perp \# \Phi)$

**unfolding** *negation-def*

```

    using measure-exchange
    by blast
  also have ... = ( $\sim \varphi \# \Gamma$ )  $\$ \vdash (\perp \# \Phi)$ 
    using ex-falso-quodlibet
      measure-tautology-right-cancel
    by blast
  finally show ?thesis .
qed

```

## 2.10 Definition of Counting Deduction

The theorem  $\Gamma \ \$ \vdash \varphi \# \Phi = \sim \varphi \# \Gamma \ \$ \vdash \perp \# \Phi$  gives rise to another kind of judgement: *how many times can a list of premises  $\Gamma$  prove a formula  $\varphi$ ?* We call this kind of judgment *counting deduction*. As with measure deduction, bits of  $\Gamma$  get "used up" with each dispatched conclusion.

```

primrec (in classical-logic)
  counting-deduction :: 'a list  $\Rightarrow$  nat  $\Rightarrow$  'a  $\Rightarrow$  bool ( $\langle \cdot \# \vdash \cdot \rangle$  [60,100,59] 60)
  where
     $\Gamma \ \# \vdash 0 \ \varphi = \text{True}$ 
  |  $\Gamma \ \# \vdash (\text{Suc } n) \ \varphi = (\exists \Psi. \text{mset } (\text{map } \text{snd } \Psi) \subseteq \# \text{mset } \Gamma \wedge$ 
       $\text{map } (\text{uncurry } (\sqcup)) \Psi \ : \vdash \varphi \wedge$ 
       $\text{map } (\text{uncurry } (\rightarrow)) \Psi \ @ \ \Gamma \ominus (\text{map } \text{snd } \Psi) \ \# \vdash n \ \varphi)$ 

```

## 2.11 Converting Back and Forth from Counting Deduction to Measure Deduction

We next show how to convert back and forth from counting deduction to measure deduction.

First, we show that trivially counting deduction is a special case of measure deduction.

```

lemma (in classical-logic) counting-deduction-to-measure-deduction:
   $\Gamma \ \# \vdash n \ \varphi = \Gamma \ \$ \vdash (\text{replicate } n \ \varphi)$ 
  by (induct n arbitrary:  $\Gamma$ , simp+)

```

We next prove a few helpful lemmas regarding counting deduction.

```

lemma (in classical-logic) counting-deduction-tautology-weaken:
  assumes  $\vdash \varphi$ 
  shows  $\Gamma \ \# \vdash n \ \varphi$ 
proof (induct n)
  case 0
  then show ?case by simp
next
  case (Suc n)
  hence  $\Gamma \ \$ \vdash (\varphi \# \text{replicate } n \ \varphi)$ 

```

```

using assms
      counting-deduction-to-measure-deduction
      measure-tautology-right-cancel
by blast
hence  $\Gamma \text{\$}\vdash \text{replicate } (\text{Suc } n) \varphi$ 
by simp
then show ?case
      using counting-deduction-to-measure-deduction
      by blast
qed

```

```

lemma (in classical-logic) counting-deduction-weaken:
  assumes  $n \leq m$ 
    and  $\Gamma \# \vdash m \varphi$ 
  shows  $\Gamma \# \vdash n \varphi$ 
proof –
  have  $\Gamma \text{\$}\vdash \text{replicate } m \varphi$ 
    using assms(2) counting-deduction-to-measure-deduction
    by blast
  hence  $\Gamma \text{\$}\vdash \text{replicate } n \varphi$ 
  by (metis append-Nil2
      assms(1)
      le-iff-add
      measure-deduction.simps(1)
      measure-deduction-generalized-witness
      replicate-add)
  thus ?thesis
    using counting-deduction-to-measure-deduction
    by blast
qed

```

```

lemma (in classical-logic) counting-deduction-implication:
  assumes  $\vdash \varphi \rightarrow \psi$ 
    and  $\Gamma \# \vdash n \varphi$ 
  shows  $\Gamma \# \vdash n \psi$ 
proof –
  have  $\text{replicate } n \psi \preceq \text{replicate } n \varphi$ 
    using stronger-theory-left-right-cons assms(1)
    by (induct n, auto)
  thus ?thesis
    using assms(2)
      measure-stronger-theory-right-antitonic
      counting-deduction-to-measure-deduction
    by blast
qed

```

Finally, we use  $\Gamma \text{\$}\vdash \varphi \# \Phi = \sim \varphi \# \Gamma \text{\$}\vdash \perp \# \Phi$  to prove that measure deduction reduces to counting deduction.

**theorem** (**in** *classical-logic*) *measure-deduction-to-counting-deduction*:

```

Γ $⊢ Φ = (∼ Φ @ Γ) #⊢ (length Φ) ⊥
proof –
have ∀ Ψ. Γ $⊢ (Φ @ Ψ) = (∼ Φ @ Γ) $⊢ (replicate (length Φ) ⊥ @ Ψ)
proof (induct Φ arbitrary: Γ)
  case Nil
  then show ?case by simp
next
case (Cons φ Φ)
  {
  fix Ψ
  have Γ $⊢ ((φ # Φ) @ Ψ) = (∼ φ # Γ) $⊢ (⊥ # Φ @ Ψ)
    using measure-negation-swap by auto
  moreover have mset (Φ @ (⊥ # Ψ)) = mset (⊥ # Φ @ Ψ)
    by simp
  ultimately have Γ $⊢ ((φ # Φ) @ Ψ) = (∼ φ # Γ) $⊢ (Φ @ (⊥ # Ψ))
    by (metis measure-msub-weaken subset-mset.order-refl)
  hence
  Γ $⊢ ((φ # Φ) @ Ψ)
    = (∼ Φ @ (∼ φ # Γ)) $⊢ (replicate (length Φ) ⊥ @ (⊥ # Ψ))
    using Cons
    by blast
  moreover have
  mset (∼ Φ @ (∼ φ # Γ)) = mset (∼ (φ # Φ) @ Γ)
  mset (replicate (length Φ) ⊥ @ (⊥ # Ψ))
    = mset (replicate (length (φ # Φ)) ⊥ @ Ψ)
    by simp+
  ultimately have
  Γ $⊢ ((φ # Φ) @ Ψ) = ∼ (φ # Φ) @ Γ $⊢ (replicate (length (φ # Φ)) ⊥
@ Ψ)
  by (metis
    append.assoc
    append-Cons
    append-Nil
    length-Cons
    replicate-append-same
    list-subtract.simps(1)
    map-ident replicate-Suc
    measure-msub-left-monotonic
    map-list-subtract-mset-containment)
  }
  then show ?case by blast
qed
thus ?thesis
  by (metis append-Nil2 counting-deduction-to-measure-deduction)
qed

```

## 2.12 Measure Deduction Soundness

The last major result for measure deduction we have to show is *soundness*. That is, judgments in measure deduction of lists of formulae can be translated into tautologies for inequalities of finitely additive probability measures over those same formulae (using the same underlying classical logic).

**lemma** (in *classical-logic*) *negated-measure-deduction*:

$$\begin{aligned} & \sim \Gamma \ \$\vdash (\varphi \# \Phi) = \\ & (\exists \Psi. \text{mset} (\text{map fst } \Psi) \subseteq\# \text{mset } \Gamma \wedge \\ & \quad \sim (\text{map} (\text{uncurry } (\backslash)) \Psi) \vdash \varphi \wedge \\ & \quad \sim (\text{map} (\text{uncurry } (\cap)) \Psi) \ @ \ \Gamma \ominus (\text{map fst } \Psi) \ \$\vdash \Phi) \end{aligned}$$

**proof** (*rule iffI*)

**assume**  $\sim \Gamma \ \$\vdash (\varphi \# \Phi)$   
**from this obtain**  $\Psi$  **where**  $\Psi$ :  
 $\text{mset} (\text{map snd } \Psi) \subseteq\# \text{mset} (\sim \Gamma)$   
 $\text{map} (\text{uncurry } (\sqcup)) \Psi \vdash \varphi$   
 $\text{map} (\text{uncurry } (\rightarrow)) \Psi \ @ \ \sim \Gamma \ominus \text{map snd } \Psi \ \$\vdash \Phi$   
**using** *measure-deduction.simps(2)*  
**by** *metis*  
**from this obtain**  $\Delta$  **where**  $\Delta$ :  
 $\text{mset } \Delta \subseteq\# \text{mset } \Gamma$   
 $\text{map snd } \Psi = \sim \Delta$   
**unfolding** *map-negation-def*  
**using** *mset-sub-map-list-exists* [**where**  $f = \sim$  **and**  $\Gamma = \Gamma$ ]  
**by** *metis*  
**let**  $? \Psi = \text{zip } \Delta (\text{map fst } \Psi)$   
**from**  $\Delta(2)$  **have**  $\text{map fst } ? \Psi = \Delta$   
**unfolding** *map-negation-def*  
**by** (*metis length-map map-fst-zip*)  
**with**  $\Delta(1)$  **have**  $\text{mset} (\text{map fst } ? \Psi) \subseteq\# \text{mset } \Gamma$   
**by** *simp*  
**moreover have**  $\forall \Delta. \text{map snd } \Psi = \sim \Delta \rightarrow$   
 $\text{map} (\text{uncurry } (\sqcup)) \Psi \preceq \sim (\text{map} (\text{uncurry } (\backslash)) (\text{zip } \Delta (\text{map fst}$

$\Psi)))$

**proof** (*induct*  $\Psi$ )

**case** *Nil*

**then show** *?case* **by** *simp*

**next**

**case** (*Cons*  $\psi \Psi$ )

**let**  $? \psi = \text{fst } \psi$

{

**fix**  $\Delta$

**assume**  $\text{map snd } (\psi \# \Psi) = \sim \Delta$

**from this obtain**  $\gamma$  **where**  $\gamma: \sim \gamma = \text{snd } \psi \ \gamma = \text{hd } \Delta$  **by** *auto*

**from**  $\langle \text{map snd } (\psi \# \Psi) = \sim \Delta \rangle$  **have**  $\text{map snd } \Psi = \sim (\text{tl } \Delta)$  **by** *auto*

**with** *Cons.hyps* **have**

$\text{map} (\text{uncurry } (\sqcup)) \Psi \preceq \sim (\text{map} (\text{uncurry } (\backslash)) (\text{zip} (\text{tl } \Delta) (\text{map fst } \Psi)))$

**by** *auto*

```

moreover
{
  fix  $\psi \ \gamma$ 
  have  $\vdash \sim(\gamma \setminus \psi) \rightarrow (\psi \sqcup \sim \gamma)$ 
    unfolding disjunction-def
      subtraction-def
      conjunction-def
      negation-def
    by (meson modus-ponens
      flip-implication
      hypothetical-syllogism)
} note tautology = this
have  $\text{uncurry } (\sqcup) = (\lambda \ \psi. (\text{fst } \psi) \sqcup (\text{snd } \psi))$ 
  by fastforce
with  $\gamma$  have  $\text{uncurry } (\sqcup) \ \psi = ?\psi \sqcup \sim \gamma$ 
  by simp
with tautology have  $\vdash \sim(\gamma \setminus ?\psi) \rightarrow \text{uncurry } (\sqcup) \ \psi$ 
  by simp
ultimately have  $\text{map } (\text{uncurry } (\sqcup)) (\psi \# \Psi) \preceq$ 
   $\sim (\text{map } (\text{uncurry } (\setminus)) ((\text{zip } ((\text{hd } \Delta) \# (\text{tl } \Delta)) (\text{map } \text{fst } (\psi \#$ 
 $\Psi))))))$ 
  using stronger-theory-left-right-cons  $\gamma(2)$ 
  by simp
hence  $\text{map } (\text{uncurry } (\sqcup)) (\psi \# \Psi) \preceq$ 
   $\sim (\text{map } (\text{uncurry } (\setminus)) (\text{zip } \Delta (\text{map } \text{fst } (\psi \# \Psi))))$ 
  using  $\langle \text{map } \text{snd } (\psi \# \Psi) = \sim \Delta \rangle$  by force
}
thus ?case by blast
qed
with  $\Psi(2) \ \Delta(2)$  have  $\sim (\text{map } (\text{uncurry } (\setminus)) \ ?\Psi) \vdash \varphi$ 
  using stronger-theory-deduction-monotonic by blast
moreover
have  $(\text{map } (\text{uncurry } (\rightarrow)) \ \Psi) @ \sim \Gamma \ominus \text{map } \text{snd } \Psi \preceq$ 
   $\sim (\text{map } (\text{uncurry } (\sqcap)) \ ?\Psi) @ \Gamma \ominus (\text{map } \text{fst } \ ?\Psi)$ 
proof –
  from  $\Delta(1)$  have  $\text{mset } (\sim \Gamma \ominus \sim \Delta) = \text{mset } (\sim (\Gamma \ominus \Delta))$ 
    by (simp add: image-mset-Diff)
  hence  $\text{mset } (\sim \Gamma \ominus \text{map } \text{snd } \Psi) = \text{mset } (\sim (\Gamma \ominus \text{map } \text{fst } \ ?\Psi))$ 
    using  $\Psi(1) \ \Delta(2) \ \langle \text{map } \text{fst } \ ?\Psi = \Delta \rangle$  by simp
  hence  $(\sim \Gamma \ominus \text{map } \text{snd } \Psi) \preceq \sim (\Gamma \ominus \text{map } \text{fst } \ ?\Psi)$ 
    by (simp add: msub-stronger-theory-intro)
  moreover have  $\forall \ \Delta. \text{map } \text{snd } \Psi = \sim \Delta \rightarrow$ 
     $\text{map } (\text{uncurry } (\rightarrow)) \ \Psi \preceq \sim (\text{map } (\text{uncurry } (\sqcap)) (\text{zip } \Delta (\text{map}$ 
 $\text{fst } \Psi)))$ 
  proof (induct  $\Psi$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\psi \ \Psi$ )

```

```

let ?ψ = fst ψ
{
  fix Δ
  assume map snd (ψ # Ψ) = ~ Δ
  from this obtain γ where γ: ~ γ = snd ψ γ = hd Δ by auto
  from ⟨map snd (ψ # Ψ) = ~ Δ⟩ have map snd Ψ = ~ (tl Δ) by auto
  with Cons.hyps have
    map (uncurry (→)) Ψ ≤ ~ (map (uncurry (∩)) (zip (tl Δ) (map fst Ψ)))
    by simp
  moreover
  {
    fix ψ γ
    have ⊢ ~ (γ ∩ ψ) → (ψ → ~ γ)
      unfolding disjunction-def
        conjunction-def
        negation-def
      by (meson modus-ponens
        flip-implication
        hypothetical-syllogism)
    } note tautology = this
  have (uncurry (→)) = (λ ψ. (fst ψ) → (snd ψ))
    by fastforce
  with γ have uncurry (→) ψ = ?ψ → ~ γ
    by simp
  with tautology have ⊢ ~ (γ ∩ ?ψ) → (uncurry (→)) ψ
    by simp
  ultimately have map (uncurry (→)) (ψ # Ψ) ≤
    ~ (map (uncurry (∩)) ((zip ((hd Δ) # (tl Δ)) (map fst (ψ #
Ψ))))))
    using stronger-theory-left-right-cons γ(2)
    by simp
  hence map (uncurry (→)) (ψ # Ψ) ≤
    ~ (map (uncurry (∩)) (zip Δ (map fst (ψ # Ψ))))
    using ⟨map snd (ψ # Ψ) = ~ Δ⟩ by force
  }
  then show ?case by blast
qed
ultimately have (map (uncurry (→)) Ψ @ ~ Γ ⊖ map snd Ψ) ≤
  (~ (map (uncurry (∩)) ?Ψ) @ ~ (Γ ⊖ (map fst ?Ψ)))
  using stronger-theory-combine Δ(2)
  by metis
thus ?thesis by simp
qed
hence ~ (map (uncurry (∩)) ?Ψ @ Γ ⊖ (map fst ?Ψ)) $⊢ Φ
  using Ψ(3) measure-stronger-theory-left-monotonic
  by blast
ultimately show ∃ Ψ. mset (map fst Ψ) ⊆# mset Γ ∧
  ~ (map (uncurry (\\)) Ψ) :⊢ φ ∧
  ~ (map (uncurry (∩)) Ψ @ Γ ⊖ (map fst Ψ)) $⊢ Φ

```

```

    by metis
next
assume  $\exists \Psi. \text{mset } (\text{map } \text{fst } \Psi) \subseteq \# \text{mset } \Gamma \wedge$ 
       $\sim (\text{map } (\text{uncurry } (\backslash)) \Psi) \vdash \varphi \wedge$ 
       $\sim (\text{map } (\text{uncurry } (\sqcap)) \Psi) @ \Gamma \ominus \text{map } \text{fst } \Psi \text{ \$}\vdash \Phi$ 
from this obtain  $\Psi$  where  $\Psi$ :
   $\text{mset } (\text{map } \text{fst } \Psi) \subseteq \# \text{mset } \Gamma$ 
   $\sim (\text{map } (\text{uncurry } (\backslash)) \Psi) \vdash \varphi$ 
   $\sim (\text{map } (\text{uncurry } (\sqcap)) \Psi) @ \Gamma \ominus \text{map } \text{fst } \Psi \text{ \$}\vdash \Phi$ 
by auto
let  $?\Psi = \text{zip } (\text{map } \text{snd } \Psi) (\sim (\text{map } \text{fst } \Psi))$ 
from  $\Psi(1)$  have  $\text{mset } (\text{map } \text{snd } ?\Psi) \subseteq \# \text{mset } (\sim \Gamma)$ 
  by (simp, metis image-mset-subseteq-mono multiset.map-comp)
moreover have  $\sim (\text{map } (\text{uncurry } (\backslash)) \Psi) \preceq \text{map } (\text{uncurry } (\sqcup)) ?\Psi$ 
proof (induct  $\Psi$ )
  case Nil
  then show ?case by simp
next
case (Cons  $\psi \Psi$ )
let  $?\gamma = \text{fst } \psi$ 
let  $?\psi = \text{snd } \psi$ 
{
  fix  $\psi \gamma$ 
  have  $\vdash (\psi \sqcup \sim \gamma) \rightarrow \sim(\gamma \backslash \psi)$ 
  unfolding disjunction-def
    subtraction-def
    conjunction-def
    negation-def
  by (meson modus-ponens
      flip-implication
      hypothetical-syllogism)
} note tautology = this
have  $\sim \circ \text{uncurry } (\backslash) = (\lambda \psi. \sim ((\text{fst } \psi) \backslash (\text{snd } \psi)))$ 
   $\text{uncurry } (\sqcup) = (\lambda (\psi, \gamma). \psi \sqcup \gamma)$ 
by fastforce+
with tautology have  $\vdash \text{uncurry } (\sqcup) (?\psi, \sim ?\gamma) \rightarrow (\sim \circ \text{uncurry } (\backslash)) \psi$ 
by fastforce
with Cons.hyps have
   $((\sim \circ \text{uncurry } (\backslash)) \psi \# \sim (\text{map } (\text{uncurry } (\backslash)) \Psi)) \preceq$ 
   $(\text{uncurry } (\sqcup) (?\psi, \sim ?\gamma) \# \text{map } (\text{uncurry } (\sqcup)) (\text{zip } (\text{map } \text{snd } \Psi) (\sim (\text{map } \text{fst } \Psi))))$ 
using stronger-theory-left-right-cons by blast
thus ?case by simp
qed
with  $\Psi(2)$  have  $\text{map } (\text{uncurry } (\sqcup)) ?\Psi \vdash \varphi$ 
using stronger-theory-deduction-monotonic by blast
moreover have  $\sim (\text{map } (\text{uncurry } (\sqcap)) \Psi) @ \Gamma \ominus \text{map } \text{fst } \Psi \preceq$ 
   $(\text{map } (\text{uncurry } (\rightarrow)) ?\Psi) @ \sim \Gamma \ominus \text{map } \text{snd } ?\Psi$ 
proof -

```



```

have  $\sim (map (uncurry (\sqcap)) \Psi) \preceq map (uncurry (\rightarrow)) ?\Psi$ 
proof (induct  $\Psi$ )
  case Nil
  then show ?case by simp
next
case (Cons  $\psi \Psi$ )
let ? $\gamma = fst \psi$ 
let ? $\psi = snd \psi$ 
{
  fix  $\psi \gamma$ 
  have  $\vdash (\psi \rightarrow \sim \gamma) \rightarrow \sim(\gamma \sqcap \psi)$ 
  unfolding disjunction-def
    conjunction-def
    negation-def
  by (meson modus-ponens
    flip-implication
    hypothetical-syllogism)
} note tautology = this
have  $\sim \circ uncurry (\sqcap) = (\lambda \psi. \sim ((fst \psi) \sqcap (snd \psi)))$ 
   $uncurry (\rightarrow) = (\lambda (\psi, \gamma). \psi \rightarrow \gamma)$ 
  by fastforce+
with tautology have  $\vdash uncurry (\rightarrow) (?\psi, \sim ?\gamma) \rightarrow (\sim \circ uncurry (\sqcap)) \psi$ 
  by fastforce
with Cons.hyps have
   $((\sim \circ uncurry (\sqcap)) \psi \# \sim (map (uncurry (\sqcap)) \Psi)) \preceq$ 
   $(uncurry (\rightarrow) (?\psi, \sim ?\gamma) \# map (uncurry (\rightarrow)) (zip (map snd \Psi) (\sim (map$ 
fst  $\Psi))))$ 
  using stronger-theory-left-right-cons by blast
then show ?case by simp
qed
moreover have  $mset (\sim (\Gamma \ominus map fst \Psi)) = mset (\sim \Gamma \ominus map snd ?\Psi)$ 
  using  $\Psi(1)$ 
  by (simp add: image-mset-Diff multiset.map-comp)
hence  $\sim (\Gamma \ominus map fst \Psi) \preceq (\sim \Gamma \ominus map snd ?\Psi)$ 
  using
    stronger-theory-reflexive
    stronger-theory-right-permutation
  by blast
ultimately show ?thesis
  using stronger-theory-combine
  by simp
qed
hence  $map (uncurry (\rightarrow)) ?\Psi @ \sim \Gamma \ominus map snd ?\Psi \ \$\vdash \Phi$ 
  using  $\Psi(3)$  measure-stronger-theory-left-monotonic by blast
ultimately show  $\sim \Gamma \ \$\vdash (\varphi \# \Phi)$ 
  using measure-deduction.simps(2) by blast
qed

```

lemma (in probability-logic) measure-deduction-soundness:

```

assumes  $\sim \Gamma \ \$\vdash \sim \Phi$ 
shows  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
proof -
have  $\forall \Gamma. \sim \Gamma \ \$\vdash \sim \Phi \longrightarrow (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
proof (induct  $\Phi$ )
  case Nil
  then show ?case
  by (simp, metis (full-types) ex-map-conv probability-non-negative sum-list-nonneg)
next
case (Cons  $\varphi \ \Phi$ )
  {
    fix  $\Gamma$ 
    assume  $\sim \Gamma \ \$\vdash \sim (\varphi \ \# \ \Phi)$ 
    hence  $\sim \Gamma \ \$\vdash (\sim \varphi \ \# \ \sim \Phi)$  by simp
    from this obtain  $\Psi$  where  $\Psi$ :
       $mset \ (map \ fst \ \Psi) \subseteq\# \ mset \ \Gamma$ 
       $\sim \ (map \ (uncurry \ (\)) \ \Psi) \ \$\vdash \sim \varphi$ 
       $\sim \ (map \ (uncurry \ (\sqcap)) \ \Psi \ @ \ \Gamma \ \ominus \ (map \ fst \ \Psi)) \ \$\vdash \sim \Phi$ 
    using negated-measure-deduction by blast
    let  $?\Gamma = \Gamma \ \ominus \ (map \ fst \ \Psi)$ 
    let  $?\Psi_1 = map \ (uncurry \ (\)) \ \Psi$ 
    let  $?\Psi_2 = map \ (uncurry \ (\sqcap)) \ \Psi$ 
    have  $(\sum \varphi' \leftarrow \Phi. \mathcal{P} \varphi') \leq (\sum \varphi \leftarrow (?\Psi_2 \ @ \ ?\Gamma). \mathcal{P} \varphi)$ 
    using Cons  $\Psi(3)$  by blast
    moreover
    have  $\mathcal{P} \varphi \leq (\sum \varphi \leftarrow ?\Psi_1. \mathcal{P} \varphi)$ 
    using  $\Psi(2)$ 
      biconditional-weaken
      list-deduction-def
      map-negation-list-implication
      set-deduction-base-theory
      implication-list-summation-inequality
    by blast
    ultimately have  $(\sum \varphi' \leftarrow (\varphi \ \# \ \Phi). \mathcal{P} \varphi') \leq (\sum \gamma \leftarrow (?\Psi_1 \ @ \ ?\Psi_2 \ @ \ ?\Gamma). \mathcal{P} \gamma)$ 
    by simp
    moreover have  $(\sum \varphi' \leftarrow (?\Psi_1 \ @ \ ?\Psi_2). \mathcal{P} \varphi') = (\sum \gamma \leftarrow (map \ fst \ \Psi). \mathcal{P} \gamma)$ 
    proof (induct  $\Psi$ )
      case Nil
      then show ?case by simp
    next
    case (Cons  $\psi \ \Psi$ )
    let  $?\Psi_1 = map \ (uncurry \ (\)) \ \Psi$ 
    let  $?\Psi_2 = map \ (uncurry \ (\sqcap)) \ \Psi$ 
    let  $? \psi_1 = uncurry \ (\) \ \psi$ 
    let  $? \psi_2 = uncurry \ (\sqcap) \ \psi$ 
    assume  $(\sum \varphi' \leftarrow (?\Psi_1 \ @ \ ?\Psi_2). \mathcal{P} \varphi') = (\sum \gamma \leftarrow (map \ fst \ \Psi). \mathcal{P} \gamma)$ 
    moreover
    {
      let  $? \gamma = fst \ \psi$ 

```

```

let ? $\psi$  = snd  $\psi$ 
have uncurry ( $\setminus$ ) = ( $\lambda$   $\psi$ . (fst  $\psi$ )  $\setminus$  (snd  $\psi$ ))
      uncurry ( $\sqcap$ ) = ( $\lambda$   $\psi$ . (fst  $\psi$ )  $\sqcap$  (snd  $\psi$ ))
      by fastforce+
moreover have  $\mathcal{P} \text{ ?}\gamma = \mathcal{P} (\text{?}\gamma \setminus \text{?}\psi) + \mathcal{P} (\text{?}\gamma \sqcap \text{?}\psi)$ 
      by (simp add: subtraction-identity)
ultimately have  $\mathcal{P} \text{ ?}\gamma = \mathcal{P} \text{ ?}\psi_1 + \mathcal{P} \text{ ?}\psi_2$ 
      by simp
}
moreover have mset ( $\text{?}\psi_1 \# \text{?}\psi_2 \# (\text{?}\Psi_1 @ \text{?}\Psi_2)$ ) =
      mset (map (uncurry ( $\setminus$ )) ( $\psi \# \Psi$ ) @ map (uncurry ( $\sqcap$ )) ( $\psi \#$ 
 $\Psi$ ))
      (is mset - = mset ?rhs)
      by simp
hence ( $\sum \varphi' \leftarrow (\text{?}\psi_1 \# \text{?}\psi_2 \# (\text{?}\Psi_1 @ \text{?}\Psi_2)). \mathcal{P} \varphi'$ ) = ( $\sum \gamma \leftarrow \text{?rhs}. \mathcal{P} \gamma$ )
      by auto
ultimately show ?case by simp
qed
moreover have mset ((map fst  $\Psi$ ) @  $\text{?}\Gamma$ ) = mset  $\Gamma$ 
      using  $\Psi(1)$ 
      by simp
hence ( $\sum \varphi' \leftarrow ((\text{map fst } \Psi) @ \text{?}\Gamma). \mathcal{P} \varphi'$ ) = ( $\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma$ )
      by (metis mset-map sum-mset-sum-list)
ultimately have ( $\sum \varphi' \leftarrow (\varphi \# \Phi). \mathcal{P} \varphi'$ )  $\leq$  ( $\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma$ )
      by simp
}
then show ?case by blast
qed
thus ?thesis using assms by blast
qed

```

## Chapter 3

# MaxSAT

We turn now to showing that counting deduction reduces to MaxSAT, the problem of finding the maximal number of satisfiable clauses in a list of clauses.

### 3.1 Definition of Relative Maximal Clause Collections

Given a list of assumptions  $\Phi$  and formula  $\varphi$ , we can think of those maximal sublists of  $\Phi$  that do not prove  $\varphi$ . While in practice we will care about  $\varphi = \perp$ , we provide a general definition in the more general axiom class *implication-logic*.

**definition** (in *implication-logic*) *relative-maximals* :: 'a list  $\Rightarrow$  'a  $\Rightarrow$  'a list set ( $\langle \mathcal{M} \rangle$ )

**where**

$$\mathcal{M} \Gamma \varphi = \{ \Phi. \text{mset } \Phi \subseteq\# \text{mset } \Gamma \wedge \neg \Phi \vdash \varphi \wedge (\forall \Psi. \text{mset } \Psi \subseteq\# \text{mset } \Gamma \longrightarrow \neg \Psi \vdash \varphi \longrightarrow \text{length } \Psi \leq \text{length } \Phi) \}$$

**lemma** (in *implication-logic*) *relative-maximals-finite*: finite ( $\mathcal{M} \Gamma \varphi$ )

**proof** –

```
{
  fix  $\Phi$ 
  assume  $\Phi \in \mathcal{M} \Gamma \varphi$ 
  hence  $\text{set } \Phi \subseteq \text{set } \Gamma$ 
     $\text{length } \Phi \leq \text{length } \Gamma$ 
  unfolding relative-maximals-def
  using mset-subset-eqD
    length-sub-mset
    mset-eq-length
  by fastforce+
```

```

}
hence  $\mathcal{M} \Gamma \varphi \subseteq \{xs. \text{set } xs \subseteq \text{set } \Gamma \wedge \text{length } xs \leq \text{length } \Gamma\}$ 
  by auto
moreover
have  $\text{finite } \{xs. \text{set } xs \subseteq \text{set } \Gamma \wedge \text{length } xs \leq \text{length } \Gamma\}$ 
  using finite-lists-length-le by blast
ultimately show ?thesis using rev-finite-subset by auto
qed

```

We know that  $\varphi$  is not a tautology if and only if the set of relative maximal sublists has an element.

**lemma** (in *implication-logic*) *relative-maximals-existence*:

$(\neg \vdash \varphi) = (\exists \Sigma. \Sigma \in \mathcal{M} \Gamma \varphi)$

**proof** (*rule iffI*)

**assume**  $\neg \vdash \varphi$

**show**  $\exists \Sigma. \Sigma \in \mathcal{M} \Gamma \varphi$

**proof** (*rule ccontr*)

**assume**  $\nexists \Sigma. \Sigma \in \mathcal{M} \Gamma \varphi$

**hence**  $\diamond: \forall \Phi. \text{mset } \Phi \subseteq\# \text{mset } \Gamma \longrightarrow$

$\neg \Phi \vdash \varphi \longrightarrow$

$(\exists \Psi. \text{mset } \Psi \subseteq\# \text{mset } \Gamma \wedge \neg \Psi \vdash \varphi \wedge \text{length } \Psi > \text{length } \Phi)$

**unfolding** *relative-maximals-def*

**by** *fastforce*

{

**fix**  $n$

**have**  $\exists \Psi. \text{mset } \Psi \subseteq\# \text{mset } \Gamma \wedge \neg \Psi \vdash \varphi \wedge \text{length } \Psi > n$

**using**  $\diamond$

**by** (*induct*  $n$ ,

*metis*

$\langle \neg \vdash \varphi \rangle$

*list.size(3)*

*list-deduction-base-theory*

*mset.simps(1)*

*subset-mset.zero-le,*

*metis*

*Nat.lessE*

*Suc-less-eq*)

}

**hence**  $\exists \Psi. \text{mset } \Psi \subseteq\# \text{mset } \Gamma \wedge \text{length } \Psi > \text{length } \Gamma$

**by** *auto*

**thus** *False*

**using** *size-mset-mono* **by** *fastforce*

**qed**

**next**

**assume**  $\exists \Sigma. \Sigma \in \mathcal{M} \Gamma \varphi$

**thus**  $\neg \vdash \varphi$

**unfolding** *relative-maximals-def*

**using** *list-deduction-weaken*

**by** *blast*

qed

**lemma** (in *implication-logic*) *relative-maximals-complement-deduction*:

**assumes**  $\Phi \in \mathcal{M} \Gamma \varphi$

**and**  $\psi \in \text{set } (\Gamma \ominus \Phi)$

**shows**  $\Phi \vdash \psi \rightarrow \varphi$

**proof** (rule *ccontr*)

**assume**  $\neg \Phi \vdash \psi \rightarrow \varphi$

**hence**  $\neg (\psi \# \Phi) \vdash \varphi$

**by** (*simp add: list-deduction-theorem*)

**moreover**

**have**  $\text{mset } \Phi \subseteq\# \text{mset } \Gamma \ \psi \in\# \text{mset } (\Gamma \ominus \Phi)$

**using** *assms*

**unfolding** *relative-maximals-def*

**by** (*blast, meson in-multiset-in-set*)

**hence**  $\text{mset } (\psi \# \Phi) \subseteq\# \text{mset } \Gamma$

**by** (*simp, metis add-mset-add-single*

*mset-subset-eq-mono-add-left-cancel*

*mset-subset-eq-single*

*subset-mset.add-diff-inverse*)

**ultimately have**  $\text{length } (\psi \# \Phi) \leq \text{length } (\Phi)$

**using** *assms*

**unfolding** *relative-maximals-def*

**by** *blast*

**thus** *False*

**by** *simp*

qed

**lemma** (in *implication-logic*) *relative-maximals-set-complement* [*simp*]:

**assumes**  $\Phi \in \mathcal{M} \Gamma \varphi$

**shows**  $\text{set } (\Gamma \ominus \Phi) = \text{set } \Gamma - \text{set } \Phi$

**proof** (rule *equalityI*)

**show**  $\text{set } (\Gamma \ominus \Phi) \subseteq \text{set } \Gamma - \text{set } \Phi$

**proof** (rule *subsetI*)

**fix**  $\psi$

**assume**  $\psi \in \text{set } (\Gamma \ominus \Phi)$

**moreover from this have**  $\Phi \vdash \psi \rightarrow \varphi$

**using** *assms*

**using** *relative-maximals-complement-deduction*

**by** *blast*

**hence**  $\psi \notin \text{set } \Phi$

**using** *assms*

*list-deduction-modus-ponens*

*list-deduction-reflection*

*relative-maximals-def*

**by** *blast*

**ultimately show**  $\psi \in \text{set } \Gamma - \text{set } \Phi$

**using** *list-subtract-set-trivial-upper-bound* [**where**  $\Gamma=\Gamma$  **and**  $\Phi=\Phi$ ]

**by** *blast*

```

qed
next
  show  $set \Gamma - set \Phi \subseteq set (\Gamma \ominus \Phi)$ 
    by (simp add: list-subtract-set-difference-lower-bound)
qed

lemma (in implication-logic) relative-maximals-complement-equiv:
  assumes  $\Phi \in \mathcal{M} \Gamma \varphi$ 
    and  $\psi \in set \Gamma$ 
  shows  $\Phi \vdash \psi \rightarrow \varphi = (\psi \notin set \Phi)$ 
proof (rule iffI)
  assume  $\Phi \vdash \psi \rightarrow \varphi$ 
  thus  $\psi \notin set \Phi$ 
    using assms(1)
      list-deduction-modus-ponens
      list-deduction-reflection
      relative-maximals-def
    by blast
next
  assume  $\psi \notin set \Phi$ 
  thus  $\Phi \vdash \psi \rightarrow \varphi$ 
    using assms relative-maximals-complement-deduction
    by auto
qed

lemma (in implication-logic) maximals-length-equiv:
  assumes  $\Phi \in \mathcal{M} \Gamma \varphi$ 
    and  $\Psi \in \mathcal{M} \Gamma \varphi$ 
  shows  $length \Phi = length \Psi$ 
  using assms
  by (simp add: dual-order.antisym relative-maximals-def)

lemma (in implication-logic) maximals-list-subtract-length-equiv:
  assumes  $\Phi \in \mathcal{M} \Gamma \varphi$ 
    and  $\Psi \in \mathcal{M} \Gamma \varphi$ 
  shows  $length (\Gamma \ominus \Phi) = length (\Gamma \ominus \Psi)$ 
proof -
  have  $length \Phi = length \Psi$ 
    using assms maximals-length-equiv
    by blast
  moreover
  have  $mset \Phi \subseteq\# mset \Gamma$ 
     $mset \Psi \subseteq\# mset \Gamma$ 
    using assms relative-maximals-def by blast+
  hence  $length (\Gamma \ominus \Phi) = length \Gamma - length \Phi$ 
     $length (\Gamma \ominus \Psi) = length \Gamma - length \Psi$ 
    by (metis list-subtract-mset-homomorphism size-Diff-submset size-mset)+
  ultimately show ?thesis by metis
qed

```

We can think of  $\Gamma \vdash \varphi$  as saying "the relative maximal sublists of  $\Gamma$  are not the entire list".

**lemma** (in *implication-logic*) *relative-maximals-max-list-deduction*:

$$\Gamma \vdash \varphi = (\forall \Phi \in \mathcal{M} \Gamma \varphi. 1 \leq \text{length} (\Gamma \ominus \Phi))$$

**proof cases**

**assume**  $\vdash \varphi$

**hence**  $\Gamma \vdash \varphi \mathcal{M} \Gamma \varphi = \{\}$

**unfolding** *relative-maximals-def*

**by** (*simp add: list-deduction-weaken*)+

**then show** *?thesis* **by** *blast*

**next**

**assume**  $\neg \vdash \varphi$

**from this obtain**  $\Omega$  **where**  $\Omega: \Omega \in \mathcal{M} \Gamma \varphi$

**using** *relative-maximals-existence* **by** *blast*

**from this have**  $\text{mset } \Omega \subseteq\# \text{mset } \Gamma$

**unfolding** *relative-maximals-def* **by** *blast*

**hence**  $\diamond: \text{length} (\Gamma \ominus \Omega) = \text{length } \Gamma - \text{length } \Omega$

**by** (*metis list-subtract-mset-homomorphism*  
*size-Diff-submset*  
*size-mset*)

**show** *?thesis*

**proof** (*cases*  $\Gamma \vdash \varphi$ )

**assume**  $\Gamma \vdash \varphi$

**from**  $\Omega$  **have**  $\text{mset } \Omega \subset\# \text{mset } \Gamma$

**by** (*metis (no-types, lifting)*

*Diff-cancel*

*Diff-eq-empty-iff*

$\langle \Gamma \vdash \varphi \rangle$

*list-deduction-monotonic*

*relative-maximals-def*

*mem-Collect-eq*

*mset-eq-setD*

*subset-mset.dual-order.not-eq-order.implies-strict*)

**hence**  $\text{length } \Omega < \text{length } \Gamma$

**using** *mset-subset-size* **by** *fastforce*

**hence**  $1 \leq \text{length } \Gamma - \text{length } \Omega$

**by** (*simp add: Suc-leI*)

**with**  $\diamond$  **have**  $1 \leq \text{length} (\Gamma \ominus \Omega)$

**by** *simp*

**with**  $\langle \Gamma \vdash \varphi \rangle \Omega$  **show** *?thesis*

**by** (*metis maximals-list-subtract-length-equiv*)

**next**

**assume**  $\neg \Gamma \vdash \varphi$

**moreover have**  $\text{mset } \Gamma \subseteq\# \text{mset } \Gamma$

**by** *simp*

**moreover have**  $\text{length } \Omega \leq \text{length } \Gamma$

**using**  $\langle \text{mset } \Omega \subseteq\# \text{mset } \Gamma \rangle$  *length-sub-mset mset-eq-length*

**by** *fastforce*

**ultimately have**  $\text{length } \Omega = \text{length } \Gamma$



```

    using  $\Omega$ 
    unfolding relative-maximals-def
    by (simp add: dual-order.antisym)
  hence  $1 > \text{length } (\Gamma \ominus \Omega)$ 
    using  $\diamond$ 
    by simp
  with  $\langle \neg \Gamma \vdash \varphi \rangle \Omega$  show ?thesis
    by fastforce
qed
qed

```

### 3.2 Definition of MaxSAT

We next turn to defining an abstract form of MaxSAT, which is largest the number of simultaneously satisfiable propositions in a list of propositions.

Unlike conventional MaxSAT, we don't actually work at the *semantic* level, i.e. constructing a model for the Tarski truth relation  $\models$ . Instead, we just count the elements in a maximal, consistent sublist (i.e., a maximal sub list  $\Sigma$  such that  $\neg \Sigma \vdash \perp$ ) of the list of assumptions  $\Gamma$  we have at hand.

Because we do not work at the semantic level, computing if  $\text{MaxSAT } \Gamma \leq n$  is not in general CoNP-Complete, as it is classically classified [1]. In the special case that the underlying logic is the *classical propositional calculus*, then the complexity is CoNP-Complete. But we could imagine the underlying logic to be linear temporal logic or even first order logic. In such cases the complexity class would be higher in the complexity hierarchy.

**definition** (in *implication-logic*) *relative-MaxSAT* :: 'a list  $\Rightarrow$  'a  $\Rightarrow$  nat ( $\langle | \cdot | \cdot \rangle$  [45])  
**where**  
 $(| \Gamma |_{\varphi}) = (\text{if } \mathcal{M} \Gamma \varphi = \{\} \text{ then } 0 \text{ else } \text{Max } \{ \text{length } \Phi \mid \Phi. \Phi \in \mathcal{M} \Gamma \varphi \})$

**abbreviation** (in *classical-logic*) *MaxSAT* :: 'a list  $\Rightarrow$  nat  
**where**  
 $\text{MaxSAT } \Gamma \equiv | \Gamma |_{\perp}$

**definition** (in *implication-logic*) *complement-relative-MaxSAT* :: 'a list  $\Rightarrow$  'a  $\Rightarrow$  nat ( $\langle || \cdot || \cdot \rangle$  [45])  
**where**  
 $(|| \Gamma ||_{\varphi}) = \text{length } \Gamma - | \Gamma |_{\varphi}$

**lemma** (in *implication-logic*) *relative-MaxSAT-intro*:  
**assumes**  $\Phi \in \mathcal{M} \Gamma \varphi$   
**shows**  $\text{length } \Phi = | \Gamma |_{\varphi}$   
**proof** –  
**have**  $\forall n \in \{ \text{length } \Psi \mid \Psi. \Psi \in \mathcal{M} \Gamma \varphi \}, n \leq \text{length } \Phi$   
 $\text{length } \Phi \in \{ \text{length } \Psi \mid \Psi. \Psi \in \mathcal{M} \Gamma \varphi \}$

**using** *assms relative-maximals-def*  
**by** *auto*  
**moreover**  
**have** *finite { length  $\Psi$  |  $\Psi. \Psi \in \mathcal{M} \Gamma \varphi$  }*  
**using** *finite-imageI relative-maximals-finite*  
**by** *simp*  
**ultimately have** *Max { length  $\Psi$  |  $\Psi. \Psi \in \mathcal{M} \Gamma \varphi$  } = length  $\Phi$*   
**using** *Max-eqI*  
**by** *blast*  
**thus** *?thesis*  
**using** *assms relative-MaxSAT-def*  
**by** *auto*  
**qed**

**lemma** (in *implication-logic*) *complement-relative-MaxSAT-intro:*

**assumes**  $\Phi \in \mathcal{M} \Gamma \varphi$   
**shows**  $\text{length} (\Gamma \ominus \Phi) = \|\Gamma\|_{\varphi}$   
**proof** –  
**have**  $\text{mset } \Phi \subseteq\# \text{mset } \Gamma$   
**using** *assms*  
**unfolding** *relative-maximals-def*  
**by** *auto*  
**moreover from this have**  $\text{length} (\Gamma \ominus \Phi) = \text{length } \Gamma - \text{length } \Phi$   
**by** (*metis list-subtract-mset-homomorphism size-Diff-submset size-mset*)  
**ultimately show** *?thesis*  
**unfolding** *complement-relative-MaxSAT-def*  
**by** (*metis assms relative-MaxSAT-intro*)  
**qed**

**lemma** (in *implication-logic*) *length-MaxSAT-decomposition:*

$\text{length } \Gamma = (\|\Gamma\|_{\varphi}) + \|\Gamma\|_{\varphi}$   
**proof** (*cases  $\mathcal{M} \Gamma \varphi = \{\}$* )  
**case** *True*  
**then show** *?thesis*  
**unfolding** *relative-MaxSAT-def*  
*complement-relative-MaxSAT-def*  
**by** *simp*  
**next**  
**case** *False*  
**from this obtain**  $\Phi$  **where**  $\Phi \in \mathcal{M} \Gamma \varphi$   
**by** *fast*  
**moreover from this have**  $\text{mset } \Phi \subseteq\# \text{mset } \Gamma$   
**unfolding** *relative-maximals-def*  
**by** *auto*  
**moreover from this have**  $\text{length} (\Gamma \ominus \Phi) = \text{length } \Gamma - \text{length } \Phi$   
**by** (*metis list-subtract-mset-homomorphism size-Diff-submset size-mset*)  
**ultimately show** *?thesis*  
**unfolding** *complement-relative-MaxSAT-def*  
**using** *list-subtract-msub-eq relative-MaxSAT-intro*

by *fastforce*  
qed

### 3.3 Reducing Counting Deduction to MaxSAT

Here we present a major result: counting deduction may be reduced to MaxSAT.

**primrec** *MaxSAT-optimal-pre-witness* :: 'a list  $\Rightarrow$  ('a list  $\times$  'a) list ( $\langle \mathfrak{W} \rangle$ )  
**where**  
 $\mathfrak{W} [] = []$   
 $|\ \mathfrak{W} (\psi \# \Psi) = (\Psi, \psi) \# \mathfrak{W} \Psi$

**lemma** *MaxSAT-optimal-pre-witness-element-inclusion*:  
 $\forall (\Delta, \delta) \in \text{set } (\mathfrak{W} \Psi). \text{set } (\mathfrak{W} \Delta) \subseteq \text{set } (\mathfrak{W} \Psi)$   
**by** (*induct*  $\Psi$ , *fastforce+*)

**lemma** *MaxSAT-optimal-pre-witness-nonelement*:

**assumes**  $\text{length } \Delta \geq \text{length } \Psi$   
**shows**  $(\Delta, \delta) \notin \text{set } (\mathfrak{W} \Psi)$   
**using** *assms*  
**proof** (*induct*  $\Psi$ )  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons*  $\psi$   $\Psi$ )  
**hence**  $\Psi \neq \Delta$  **by** *auto*  
**then show** *?case* **using** *Cons* **by** *simp*  
**qed**

**lemma** *MaxSAT-optimal-pre-witness-distinct*: *distinct* ( $\mathfrak{W} \Psi$ )  
**by** (*induct*  $\Psi$ , *simp*, *simp add: MaxSAT-optimal-pre-witness-nonelement*)

**lemma** *MaxSAT-optimal-pre-witness-length-iff-eq*:

$\forall (\Delta, \delta) \in \text{set } (\mathfrak{W} \Psi). \forall (\Sigma, \sigma) \in \text{set } (\mathfrak{W} \Psi). (\text{length } \Delta = \text{length } \Sigma) = ((\Delta, \delta) = (\Sigma, \sigma))$   
**proof** (*induct*  $\Psi$ )  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons*  $\psi$   $\Psi$ )  
**{**  
**fix**  $\Delta$   
**fix**  $\delta$   
**assume**  $(\Delta, \delta) \in \text{set } (\mathfrak{W} (\psi \# \Psi))$   
**and**  $\text{length } \Delta = \text{length } \Psi$   
**hence**  $(\Delta, \delta) = (\Psi, \psi)$   
**by** (*simp add: MaxSAT-optimal-pre-witness-nonelement*)  
**}**

**hence**  $\forall (\Delta, \delta) \in \text{set } (\mathfrak{V} (\psi \# \Psi)). (\text{length } \Delta = \text{length } \Psi) = ((\Delta, \delta) = (\Psi, \psi))$   
**by** *blast*  
**with** *Cons* **show** *?case*  
**by** *auto*  
**qed**

**lemma** *mset-distinct-msub-down*:  
**assumes**  $mset A \subseteq\# mset B$   
**and** *distinct B*  
**shows** *distinct A*  
**using** *assms*  
**by** (*meson distinct-append mset-le-perm-append perm-distinct-iff*)

**lemma** *mset-remdups-set-sub-iff*:  
 $(mset (\text{remdups } A) \subseteq\# mset (\text{remdups } B)) = (\text{set } A \subseteq \text{set } B)$   
**proof** –  
**have**  $\forall B. (mset (\text{remdups } A) \subseteq\# mset (\text{remdups } B)) = (\text{set } A \subseteq \text{set } B)$   
**proof** (*induct A*)  
**case** *Nil*  
**then show** *?case* **by** *simp*  
**next**  
**case** (*Cons a A*)  
**then show** *?case*  
**proof** (*cases a ∈ set A*)  
**case** *True*  
**then show** *?thesis* **using** *Cons* **by** *auto*  
**next**  
**case** *False*  
**{**  
**fix** *B*  
**have**  $(mset (\text{remdups } (a \# A)) \subseteq\# mset (\text{remdups } B)) = (\text{set } (a \# A) \subseteq \text{set } B)$   
**proof** (*rule iffI*)  
**assume** *assm: mset (remdups (a # A)) ⊆# mset (remdups B)*  
**hence**  $mset (\text{remdups } A) \subseteq\# mset (\text{remdups } B) - \{\#a\#$   
**using** *False*  
**by** (*simp add: insert-subset-eq-iff*)  
**hence**  $mset (\text{remdups } A) \subseteq\# mset (\text{remdups } (\text{removeAll } a B))$   
**by** (*metis diff-subset-eq-self*  
*distinct-remdups*  
*distinct-remove1-removeAll*  
*mset-distinct-msub-down*  
*mset-remove1*  
*set-eq-iff-mset-eq-distinct*  
*set-remdups set-removeAll*)  
**hence**  $\text{set } A \subseteq \text{set } (\text{removeAll } a B)$   
**using** *Cons.hyps* **by** *blast*  
**moreover from** *assm False* **have**  $a \in \text{set } B$   
**using** *mset-subset-eq-insertD* **by** *fastforce*

```

ultimately show  $set (a \# A) \subseteq set B$ 
  by auto
next
assume  $assm: set (a \# A) \subseteq set B$ 
hence  $set A \subseteq set (removeAll a B)$  using False
  by auto
hence  $mset (remdups A) \subseteq\# mset (remdups B) - \{\#a\}$ 
  by (metis Cons.hyps
        distinct-remdups
        mset-remdups-subset-eq
        mset-remove1 remove-code(1)
        set-remdups set-remove1-eq
        set-removeAll
        subset-mset.dual-order.trans)
moreover from  $assm$  False have  $a \in set B$  by auto
ultimately show  $mset (remdups (a \# A)) \subseteq\# mset (remdups B)$ 
  by (simp add: False insert-subset-eq-iff)
qed
}
then show ?thesis by simp
qed
qed
thus ?thesis by blast
qed

lemma range-characterization:
  ( $mset X = mset [0..<length X]$ ) = ( $distinct X \wedge (\forall x \in set X. x < length X)$ )
proof (rule iffI)
  assume  $mset X = mset [0..<length X]$ 
  thus  $distinct X \wedge (\forall x \in set X. x < length X)$ 
  by (metis atLeastLessThan-iff count-mset-0-iff distinct-count-atmost-1 distinct-upt
set-upt)
next
assume  $distinct X \wedge (\forall x \in set X. x < length X)$ 
moreover
{
  fix  $n$ 
  have  $\forall X. n = length X \longrightarrow$ 
     $distinct X \wedge (\forall x \in set X. x < length X) \longrightarrow$ 
     $mset X = mset [0..<length X]$ 
proof (induct  $n$ )
  case 0
  then show ?case by simp
next
  case (Suc  $n$ )
  {
    fix  $X$ 
    assume  $A: n + 1 = length X$ 
    and  $B: distinct X$ 

```

```

    and C:  $\forall x \in \text{set } X. x < \text{length } X$ 
  have n  $\in \text{set } X$ 
  proof (rule ccontr)
    assume n  $\notin \text{set } X$ 
    from A have A':  $n = \text{length } (\text{tl } X)$ 
    by simp
    from B have B': distinct (tl X)
    by (simp add: distinct-tl)
    have C':  $\forall x \in \text{set } (\text{tl } X). x < \text{length } (\text{tl } X)$ 
    by (metis
        A
        A'
        C
         $\langle n \notin \text{set } X \rangle$ 
        Suc-eq-plus1
        Suc-le-eq
        Suc-le-mono
        le-less
        list.set-sel(2)
        list.size(3)
        nat.simps(3))
    from A' B' C' Suc have  $\text{mset } (\text{tl } X) = \text{mset } [0..<n]$ 
    by blast
    from A have  $X = \text{hd } X \# \text{tl } X$ 
    by (metis Suc-eq-plus1 list.exhaust-sel list.size(3) nat.simps(3))
    with B  $\langle \text{mset } (\text{tl } X) = \text{mset } [0..<n] \rangle$  have  $\text{hd } X \notin \text{set } [0..<n]$ 
    by (metis distinct.simps(2) mset-eq-setD)
    hence  $\text{hd } X \geq n$  by simp
    with C  $\langle n \notin \text{set } X \rangle$   $\langle X = \text{hd } X \# \text{tl } X \rangle$  show False
    by (metis A Suc-eq-plus1 Suc-le-eq le-neq-trans list.set-intros(1) not-less)
  qed
  let ?X' = remove1 n X
  have A':  $n = \text{length } ?X'$ 
  by (metis A  $\langle n \in \text{set } X \rangle$  diff-add-inverse2 length-remove1)
  have B': distinct ?X'
  by (simp add: B)
  have C':  $\forall x \in \text{set } ?X'. x < \text{length } ?X'$ 
  by (metis A A' B C
      DiffE
      Suc-eq-plus1
      Suc-le-eq
      Suc-le-mono
      le-neq-trans
      set-remove1-eq
      singletonI)
  hence  $\text{mset } ?X' = \text{mset } [0..<n]$ 
  using A' B' C' Suc
  by auto
  hence  $\text{mset } (n \# ?X') = \text{mset } [0..<n+1]$ 

```

```

    by simp
  hence mset X = mset [0.. $\text{length } X$ ]
    by (metis A ⟨ $n \in \text{set } X$ ⟩ perm-remove)
  }
  then show ?case by fastforce
qed
}
ultimately show mset X = mset [0.. $\text{length } X$ ]
  by blast
qed

```

**lemma** *distinct-pigeon-hole*:

```

  fixes X :: nat list
  assumes distinct X
    and X ≠ []
  shows ∃ n ∈ set X. n + 1 ≥ length X
proof (rule ccontr)
  assume *: ¬ (∃ n ∈ set X. length X ≤ n + 1)
  hence ∀ n ∈ set X. n < length X by fastforce
  hence mset X = mset [0.. $\text{length } X$ ]
    using assms(1) range-characterization
    by fastforce
  with assms(2) have length X - 1 ∈ set X
    by (metis
      diff-zero
      last-in-set
      last-upt
      length-greater-0-conv
      length-upt mset-eq-setD)
  with * show False
    by (metis One-nat-def Suc-eq-plus1 Suc-pred le-refl length-pos-if-in-set)
qed

```

**lemma** *MaxSAT-optimal-pre-witness-pigeon-hole*:

```

  assumes mset Σ ⊆# mset (⋈ Ψ)
    and Σ ≠ []
  shows ∃ (Δ, δ) ∈ set Σ. length Δ + 1 ≥ length Σ
proof -
  have distinct Σ
    using assms
      MaxSAT-optimal-pre-witness-distinct
      mset-distinct-msub-down
    by blast
  with assms(1) have distinct (map (length ∘ fst) Σ)
proof (induct Σ)
  case Nil
    then show ?case by simp
  next
  case (Cons σ Σ)

```

**hence**  $mset \Sigma \subseteq\# mset (\mathfrak{W} \Psi)$   
*distinct*  $\Sigma$   
**by** (*metis* *mset.simps(2)* *mset-subset-eq-insertD* *subset-mset-def*, *simp*)  
**with** *Cons.hyps* **have** *distinct* (*map* ( $\lambda a. length (fst a)$ )  $\Sigma$ ) **by** *simp*  
**moreover**  
**obtain**  $\delta \Delta$  **where**  $\sigma = (\Delta, \delta)$   
**by** *fastforce*  
**hence**  $(\Delta, \delta) \in set (\mathfrak{W} \Psi)$   
**using** *Cons.prem*s *mset-subset-eq-insertD*  
**by** *fastforce*  
**hence**  $\forall (\Sigma, \sigma) \in set (\mathfrak{W} \Psi). (length \Delta = length \Sigma) = ((\Delta, \delta) = (\Sigma, \sigma))$   
**using** *MaxSAT-optimal-pre-witness-length-iff-eq* [**where**  $\Psi = \Psi$ ]  
**by** *fastforce*  
**hence**  $\forall (\Sigma, \sigma) \in set \Sigma. (length \Delta = length \Sigma) = ((\Delta, \delta) = (\Sigma, \sigma))$   
**using**  $\langle mset \Sigma \subseteq\# mset (\mathfrak{W} \Psi) \rangle$   
**by** (*metis* (*no-types*, *lifting*) *Un-iff mset-le-perm-append perm-set-eq set-append*)  
**hence**  $length (fst \sigma) \notin set (map (\lambda a. length (fst a)) \Sigma)$   
**using** *Cons.prem*s(2)  $\langle \sigma = (\Delta, \delta) \rangle$   
**by** *fastforce*  
**ultimately show** *?case* **by** *simp*  
**qed**  
**moreover have**  $length (map (length \circ fst) \Sigma) = length \Sigma$  **by** *simp*  
**moreover have**  $map (length \circ fst) \Sigma \neq []$  **using** *assms* **by** *simp*  
**ultimately show** *?thesis*  
**using** *distinct-pigeon-hole*  
**by** *fastforce*  
**qed**

**abbreviation (in classical-logic)**  
*MaxSAT-optimal-witness*  $:: 'a \Rightarrow 'a list \Rightarrow ('a \times 'a) list (\langle \mathfrak{W} \rangle)$   
**where**  $\mathfrak{W} \varphi \Xi \equiv map (\lambda (\Psi, \psi). (\Psi \rightarrow \varphi, \psi)) (\mathfrak{W} \Xi)$

**abbreviation (in classical-logic)**  
*disjunction-MaxSAT-optimal-witness*  $:: 'a \Rightarrow 'a list \Rightarrow 'a list (\langle \mathfrak{W}_{\sqcup} \rangle)$   
**where**  $\mathfrak{W}_{\sqcup} \varphi \Psi \equiv map (uncurry (\sqcup)) (\mathfrak{W} \varphi \Psi)$

**abbreviation (in classical-logic)**  
*implication-MaxSAT-optimal-witness*  $:: 'a \Rightarrow 'a list \Rightarrow 'a list (\langle \mathfrak{W}_{\rightarrow} \rangle)$   
**where**  $\mathfrak{W}_{\rightarrow} \varphi \Psi \equiv map (uncurry (\rightarrow)) (\mathfrak{W} \varphi \Psi)$

**lemma (in classical-logic)** *MaxSAT-optimal-witness-conjunction-identity*:

$\vdash \sqcap (\mathfrak{W}_{\sqcup} \varphi \Psi) \leftrightarrow (\varphi \sqcup \sqcap \Psi)$

**proof** (*induct*  $\Psi$ )

**case** *Nil*

**then show** *?case*

**unfolding** *biconditional-def*

*disjunction-def*

**using** *axiom-k*

*modus-ponens*



```

      verum-tautology
    by (simp, blast)
next
case (Cons  $\psi$   $\Psi$ )
have  $\vdash (\Psi \rightarrow \varphi) \leftrightarrow (\bigwedge \Psi \rightarrow \varphi)$ 
  by (simp add: list-curry-uncurry)
hence  $\vdash \bigwedge (\text{map } (\text{uncurry } \sqcup) (\mathfrak{M} \varphi (\psi \# \Psi)))$ 
   $\leftrightarrow ((\bigwedge \Psi \rightarrow \varphi \sqcup \psi) \sqcap \bigwedge (\text{map } (\text{uncurry } \sqcup) (\mathfrak{M} \varphi \Psi)))$ 
  unfolding biconditional-def
  using conjunction-monotonic
  disjunction-monotonic
  by simp
moreover have  $\vdash ((\bigwedge \Psi \rightarrow \varphi \sqcup \psi) \sqcap \bigwedge (\text{map } (\text{uncurry } \sqcup) (\mathfrak{M} \varphi \Psi)))$ 
   $\leftrightarrow ((\bigwedge \Psi \rightarrow \varphi \sqcup \psi) \sqcap (\varphi \sqcup \bigwedge \Psi))$ 
  using Cons.hyps biconditional-conjunction-weaken-rule
  by blast
moreover
{
  fix  $\varphi \psi \chi$ 
  have  $\vdash ((\chi \rightarrow \varphi \sqcup \psi) \sqcap (\varphi \sqcup \chi)) \leftrightarrow (\varphi \sqcup (\psi \sqcap \chi))$ 
  proof -
    let  $? \varphi = ((\langle \chi \rangle \rightarrow \langle \varphi \rangle \sqcup \langle \psi \rangle) \sqcap (\langle \varphi \rangle \sqcup \langle \chi \rangle)) \leftrightarrow (\langle \varphi \rangle \sqcup (\langle \psi \rangle \sqcap \langle \chi \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  by fastforce
    hence  $\vdash (\langle ? \varphi \rangle)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
}
ultimately have  $\vdash \bigwedge (\text{map } (\text{uncurry } \sqcup) (\mathfrak{M} \varphi (\psi \# \Psi))) \leftrightarrow (\varphi \sqcup (\psi \sqcap \bigwedge \Psi))$ 
  using biconditional-transitivity-rule
  by blast
then show ?case by simp
qed

```

**lemma (in classical-logic) MaxSAT-optimal-witness-deduction:**

```

 $\vdash \mathfrak{M}_{\sqcup} \varphi \Psi \rightarrow \varphi \leftrightarrow \Psi \rightarrow \varphi$ 
proof -
have  $\vdash \mathfrak{M}_{\sqcup} \varphi \Psi \rightarrow \varphi \leftrightarrow (\bigwedge (\mathfrak{M}_{\sqcup} \varphi \Psi) \rightarrow \varphi)$ 
  by (simp add: list-curry-uncurry)
moreover
{
  fix  $\alpha \beta \gamma$ 
  have  $\vdash (\alpha \leftrightarrow \beta) \rightarrow ((\alpha \rightarrow \gamma) \leftrightarrow (\beta \rightarrow \gamma))$ 
  proof -
    let  $? \varphi = (\langle \alpha \rangle \leftrightarrow \langle \beta \rangle) \rightarrow ((\langle \alpha \rangle \rightarrow \langle \gamma \rangle) \leftrightarrow (\langle \beta \rangle \rightarrow \langle \gamma \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  by fastforce
    hence  $\vdash (\langle ? \varphi \rangle)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
}
qed

```

```

}
ultimately have  $\vdash \mathfrak{W}_{\sqcup} \varphi \Psi \text{ :} \rightarrow \varphi \leftrightarrow ((\varphi \sqcup \prod \Psi) \rightarrow \varphi)$ 
  using modus-ponens
         biconditional-transitivity-rule
         MaxSAT-optimal-witness-conjunction-identity
  by blast
moreover
{
  fix  $\alpha \beta$ 
  have  $\vdash ((\alpha \sqcup \beta) \rightarrow \alpha) \leftrightarrow (\beta \rightarrow \alpha)$ 
  proof -
    let  $? \varphi = (((\langle \alpha \rangle \sqcup \langle \beta \rangle) \rightarrow \langle \alpha \rangle) \leftrightarrow (\langle \beta \rangle \rightarrow \langle \alpha \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  by fastforce
    hence  $\vdash (| ? \varphi |)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
}
ultimately have  $\vdash \mathfrak{W}_{\sqcup} \varphi \Psi \text{ :} \rightarrow \varphi \leftrightarrow (\prod \Psi \rightarrow \varphi)$ 
  using biconditional-transitivity-rule by blast
  thus ?thesis
  using biconditional-symmetry-rule
         biconditional-transitivity-rule
         list-curry-uncurry
  by blast
qed

lemma (in classical-logic) optimal-witness-split-identity:
 $\vdash (\mathfrak{W}_{\sqcup} \varphi (\psi \# \Xi)) \text{ :} \rightarrow \varphi \rightarrow (\mathfrak{W}_{\rightarrow} \varphi (\psi \# \Xi)) \text{ :} \rightarrow \varphi \rightarrow \Xi \text{ :} \rightarrow \varphi$ 
proof (induct  $\Xi$ )
  case Nil
  have  $\vdash ((\varphi \sqcup \psi) \rightarrow \varphi) \rightarrow ((\varphi \rightarrow \psi) \rightarrow \varphi) \rightarrow \varphi$ 
  proof -
    let  $? \varphi = (((\langle \varphi \rangle \sqcup \langle \psi \rangle) \rightarrow \langle \varphi \rangle) \rightarrow ((\langle \varphi \rangle \rightarrow \langle \psi \rangle) \rightarrow \langle \varphi \rangle) \rightarrow \langle \varphi \rangle)$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  by fastforce
    hence  $\vdash (| ? \varphi |)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
  then show ?case by simp
next
  case (Cons  $\xi \Xi$ )
  let  $?A = \mathfrak{W}_{\sqcup} \varphi \Xi \text{ :} \rightarrow \varphi$ 
  let  $?B = \mathfrak{W}_{\rightarrow} \varphi \Xi \text{ :} \rightarrow \varphi$ 
  let  $?X = \Xi \text{ :} \rightarrow \varphi$ 
  from Cons.hyps have  $\vdash ((?X \sqcup \psi) \rightarrow ?A) \rightarrow ((?X \rightarrow \psi) \rightarrow ?B) \rightarrow ?X$  by simp
  moreover
  have  $\vdash (((?X \sqcup \psi) \rightarrow ?A) \rightarrow ((?X \rightarrow \psi) \rightarrow ?B) \rightarrow ?X)$ 
     $\rightarrow ((\xi \rightarrow ?X \sqcup \psi) \rightarrow (?X \sqcup \xi) \rightarrow ?A) \rightarrow (((\xi \rightarrow ?X) \rightarrow \psi) \rightarrow (?X \rightarrow \xi)$ 
     $\rightarrow ?B) \rightarrow \xi \rightarrow ?X$ 
  proof -

```

**let**  $?φ = (((\langle ?X \rangle \sqcup \langle \psi \rangle) \rightarrow \langle ?A \rangle) \rightarrow ((\langle ?X \rangle \rightarrow \langle \psi \rangle) \rightarrow \langle ?B \rangle) \rightarrow \langle ?X \rangle \rightarrow$   
 $((\langle \xi \rangle \rightarrow \langle ?X \rangle \sqcup \langle \psi \rangle) \rightarrow (\langle ?X \rangle \sqcup \langle \xi \rangle) \rightarrow \langle ?A \rangle) \rightarrow$   
 $((\langle \xi \rangle \rightarrow \langle ?X \rangle) \rightarrow \langle \psi \rangle) \rightarrow (\langle ?X \rangle \rightarrow \langle \xi \rangle) \rightarrow \langle ?B \rangle) \rightarrow$   
 $\langle \xi \rangle \rightarrow$   
 $\langle ?X \rangle$   
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?φ$  **by** *fastforce*  
**hence**  $\vdash (\langle ?φ \rangle)$  **using** *propositional-semantic* **by** *blast*  
**thus** *?thesis* **by** *simp*  
**qed**  
**ultimately**  
**have**  $\vdash ((\xi \rightarrow ?X \sqcup \psi) \rightarrow (?X \sqcup \xi) \rightarrow ?A) \rightarrow (((\xi \rightarrow ?X) \rightarrow \psi) \rightarrow (?X \rightarrow \xi)$   
 $\rightarrow ?B) \rightarrow \xi \rightarrow ?X$   
**using** *modus-ponens*  
**by** *blast*  
**thus** *?case* **by** *simp*  
**qed**

**lemma** (in *classical-logic*) *disj-conj-impl-duality*:  
 $\vdash (\varphi \rightarrow \chi \sqcap \psi \rightarrow \chi) \leftrightarrow ((\varphi \sqcup \psi) \rightarrow \chi)$   
**proof** –  
**let**  $?φ = ((\langle \varphi \rangle \rightarrow \langle \chi \rangle \sqcap \langle \psi \rangle \rightarrow \langle \chi \rangle) \leftrightarrow ((\langle \varphi \rangle \sqcup \langle \psi \rangle) \rightarrow \langle \chi \rangle))$   
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?φ$  **by** *fastforce*  
**hence**  $\vdash (\langle ?φ \rangle)$  **using** *propositional-semantic* **by** *blast*  
**thus** *?thesis* **by** *simp*  
**qed**

**lemma** (in *classical-logic*) *weak-disj-of-conj-equiv*:  
 $(\forall \sigma \in set \Sigma. \sigma \vdash \varphi) = \vdash \sqcup (map \sqcap \Sigma) \rightarrow \varphi$   
**proof** (*induct*  $\Sigma$ )  
**case** *Nil*  
**then show** *?case*  
**by** (*simp add: ex-falso-quodlibet*)  
**next**  
**case** (*Cons*  $\sigma \Sigma$ )  
**have**  $(\forall \sigma' \in set (\sigma \# \Sigma). \sigma' \vdash \varphi) = (\sigma \vdash \varphi \wedge (\forall \sigma' \in set \Sigma. \sigma' \vdash \varphi))$  **by** *simp*  
**also have**  $\dots = (\vdash \sigma \rightarrow \varphi \wedge \vdash \sqcup (map \sqcap \Sigma) \rightarrow \varphi)$  **using** *Cons.hyps list-deduction-def*  
**by** *simp*  
**also have**  $\dots = (\vdash \sqcap \sigma \rightarrow \varphi \wedge \vdash \sqcup (map \sqcap \Sigma) \rightarrow \varphi)$   
**using** *list-curry-uncurry weak-biconditional-weaken* **by** *blast*  
**also have**  $\dots = (\vdash \sqcap \sigma \rightarrow \varphi \sqcap \sqcup (map \sqcap \Sigma) \rightarrow \varphi)$  **by** *simp*  
**also have**  $\dots = (\vdash (\sqcap \sigma \sqcup \sqcup (map \sqcap \Sigma)) \rightarrow \varphi)$   
**using** *disj-conj-impl-duality weak-biconditional-weaken* **by** *blast*  
**finally show** *?case* **by** *simp*  
**qed**

**lemma** (in *classical-logic*) *arbitrary-disj-concat-equiv*:  
 $\vdash \sqcup (\Phi @ \Psi) \leftrightarrow (\sqcup \Phi \sqcup \sqcup \Psi)$   
**proof** (*induct*  $\Phi$ )  
**case** *Nil*

```

then show ?case
  by (simp,
    meson ex-falso-quodlibet
      modus-ponens
      biconditional-introduction
      disjunction-elimination
      disjunction-right-introduction
      trivial-implication)
next
  case (Cons  $\varphi$   $\Phi$ )
  have  $\vdash \sqcup (\Phi @ \Psi) \leftrightarrow (\sqcup \Phi \sqcup \sqcup \Psi) \rightarrow (\varphi \sqcup \sqcup (\Phi @ \Psi)) \leftrightarrow ((\varphi \sqcup \sqcup \Phi) \sqcup \sqcup \Psi)$ 
  proof -
    let ? $\varphi$  =
       $(\langle \sqcup (\Phi @ \Psi) \rangle \leftrightarrow (\langle \sqcup \Phi \rangle \sqcup \langle \sqcup \Psi \rangle)) \rightarrow (\langle \varphi \rangle \sqcup \langle \sqcup (\Phi @ \Psi) \rangle) \leftrightarrow ((\langle \varphi \rangle \sqcup \langle \sqcup \Phi \rangle) \sqcup \langle \sqcup \Psi \rangle)$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\langle ?\varphi \rangle)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
  then show ?case using Cons modus-ponens by simp
qed

lemma (in classical-logic) arbitrary-conj-concat-equiv:
   $\vdash \sqcap (\Phi @ \Psi) \leftrightarrow (\sqcap \Phi \sqcap \sqcap \Psi)$ 
proof (induct  $\Phi$ )
  case Nil
  then show ?case
  by (simp,
    meson modus-ponens
      biconditional-introduction
      conjunction-introduction
      conjunction-right-elimination
      verum-tautology)
next
  case (Cons  $\varphi$   $\Phi$ )
  have  $\vdash \sqcap (\Phi @ \Psi) \leftrightarrow (\sqcap \Phi \sqcap \sqcap \Psi) \rightarrow (\varphi \sqcap \sqcap (\Phi @ \Psi)) \leftrightarrow ((\varphi \sqcap \sqcap \Phi) \sqcap \sqcap \Psi)$ 
  proof -
    let ? $\varphi$  =
       $(\langle \sqcap (\Phi @ \Psi) \rangle \leftrightarrow (\langle \sqcap \Phi \rangle \sqcap \langle \sqcap \Psi \rangle)) \rightarrow (\langle \varphi \rangle \sqcap \langle \sqcap (\Phi @ \Psi) \rangle) \leftrightarrow ((\langle \varphi \rangle \sqcap \langle \sqcap \Phi \rangle) \sqcap \langle \sqcap \Psi \rangle)$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\langle ?\varphi \rangle)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed
  then show ?case using Cons modus-ponens by simp
qed

```

```

lemma (in classical-logic) conj-absorption:
  assumes  $\chi \in \text{set } \Phi$ 
  shows  $\vdash \sqcap \Phi \leftrightarrow (\chi \sqcap \sqcap \Phi)$ 
  using assms
proof (induct  $\Phi$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\varphi \Phi$ )
  then show ?case
  proof (cases  $\varphi = \chi$ )
    case True
    then show ?thesis
    by (simp,
      metis biconditional-def
      implication-distribution
      trivial-implication
      weak-biconditional-weaken
      weak-conjunction-deduction-equivalence)
  next
  case False
  then show ?thesis
  by (metis Cons.prem
    arbitrary-conjunction.simps(2)
    modus-ponens
    arbitrary-conjunction-antitone
    biconditional-introduction
    remdups.simps(2)
    set-remdups
    set-subset-Cons)

  qed
qed

lemma (in classical-logic) conj-extract:  $\vdash \sqcup (\text{map } ((\sqcap) \varphi) \Psi) \leftrightarrow (\varphi \sqcap \sqcup \Psi)$ 
proof (induct  $\Psi$ )
  case Nil
  then show ?case
  by (simp add: ex-falso-quodlibet biconditional-def conjunction-right-elimination)
next
  case (Cons  $\psi \Psi$ )
  have  $\vdash \sqcup (\text{map } ((\sqcap) \varphi) \Psi) \leftrightarrow (\varphi \sqcap \sqcup \Psi)$ 
     $\rightarrow ((\varphi \sqcap \psi) \sqcup \sqcup (\text{map } ((\sqcap) \varphi) \Psi)) \leftrightarrow (\varphi \sqcap (\psi \sqcup \sqcup \Psi))$ 
  proof –
    let  $?\varphi = \langle \sqcup (\text{map } ((\sqcap) \varphi) \Psi) \rangle \leftrightarrow (\langle \varphi \rangle \sqcap \langle \sqcup \Psi \rangle)$ 
       $\rightarrow ((\langle \varphi \rangle \sqcap \langle \psi \rangle) \sqcup \langle \sqcup (\text{map } ((\sqcap) \varphi) \Psi) \rangle) \leftrightarrow (\langle \varphi \rangle \sqcap (\langle \psi \rangle \sqcup \langle \sqcup \Psi \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\langle ?\varphi \rangle)$  using propositional-semantic by blast
    thus ?thesis by simp
  qed

```

**then show** *?case* **using** *Cons modus-ponens* **by** *simp*  
**qed**

**lemma** (in *classical-logic*) *conj-multi-extract*:

$\vdash \sqcup (map \sqcap (map ((@) \Delta) \Sigma)) \leftrightarrow (\sqcap \Delta \sqcap \sqcup (map \sqcap \Sigma))$

**proof** (*induct*  $\Sigma$ )

**case** *Nil*

**then show** *?case*

**by** (*simp*, *metis list.simps(8) arbitrary-disjunction.simps(1) conj-extract*)

**next**

**case** (*Cons*  $\sigma$   $\Sigma$ )

**moreover have**

$\vdash \sqcup (map \sqcap (map ((@) \Delta) \Sigma)) \leftrightarrow (\sqcap \Delta \sqcap \sqcup (map \sqcap \Sigma))$

$\rightarrow \sqcap (\Delta @ \sigma) \leftrightarrow (\sqcap \Delta \sqcap \sqcap \sigma)$

$\rightarrow (\sqcap (\Delta @ \sigma) \sqcup \sqcup (map (\sqcap \circ (@) \Delta) \Sigma)) \leftrightarrow (\sqcap \Delta \sqcap (\sqcap \sigma \sqcup \sqcup (map \sqcap \Sigma)))$

**proof** –

**let** *? $\varphi$*  =

$\langle \sqcup (map \sqcap (map ((@) \Delta) \Sigma)) \rangle \leftrightarrow (\langle \sqcap \Delta \rangle \sqcap \langle \sqcup (map \sqcap \Sigma) \rangle)$

$\rightarrow \langle \sqcap (\Delta @ \sigma) \rangle \leftrightarrow (\langle \sqcap \Delta \rangle \sqcap \langle \sqcap \sigma \rangle)$

$\rightarrow (\langle \sqcap (\Delta @ \sigma) \rangle \sqcup \langle \sqcup (map (\sqcap \circ (@) \Delta) \Sigma) \rangle) \leftrightarrow (\langle \sqcap \Delta \rangle \sqcap (\langle \sqcap \sigma \rangle \sqcup \langle \sqcup (map \sqcap \Sigma) \rangle))$

**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  **by** *fastforce*

**hence**  $\vdash (\ ?\varphi \ )$  **using** *propositional-semantic* **by** *blast*

**thus** *?thesis* **by** *simp*

**qed**

**hence**

$\vdash (\sqcap (\Delta @ \sigma) \sqcup \sqcup (map (\sqcap \circ (@) \Delta) \Sigma)) \leftrightarrow (\sqcap \Delta \sqcap (\sqcap \sigma \sqcup \sqcup (map \sqcap \Sigma)))$

**using** *Cons.hyps arbitrary-conj-concat-equiv modus-ponens* **by** *blast*

**then show** *?case* **by** *simp*

**qed**

**lemma** (in *classical-logic*) *extract-inner-concat*:

$\vdash \sqcup (map (\sqcap \circ (map \text{snd} \circ (@) \Delta)) \Psi) \leftrightarrow (\sqcap (map \text{snd} \Delta) \sqcap \sqcup (map (\sqcap \circ map \text{snd}) \Psi))$

**proof** (*induct*  $\Delta$ )

**case** *Nil*

**then show** *?case*

**by** (*simp*,

*meson modus-ponens*

*biconditional-introduction*

*conjunction-introduction*

*conjunction-right-elimination*

*verum-tautology*)

**next**

**case** (*Cons*  $\chi$   $\Delta$ )

**let** *? $\Delta'$*  = *map snd*  $\Delta$

**let** *? $\chi'$*  = *snd*  $\chi$

**let**  $?\Pi = \lambda\varphi. \sqcap (\text{map } \text{snd } \varphi)$   
**let**  $?\Pi\Delta = \lambda\varphi. \sqcap (?\Delta' @ \text{map } \text{snd } \varphi)$   
**from** *Cons* **have**  
 $\vdash \sqcup (\text{map } ?\Pi\Delta \Psi) \leftrightarrow (\sqcap ?\Delta' \sqcap \sqcup (\text{map } ?\Pi \Psi))$   
**by** *auto*  
**moreover have**  $\star: \text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) = \text{map } ((\sqcap ?\chi') \circ \text{map } ?\Pi\Delta)$   
**by** *fastforce*  
**have**  $\sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) \Psi) = \sqcup (\text{map } ((\sqcap ?\chi') (\text{map } ?\Pi\Delta \Psi))$   
**by** (*simp add: \**)  
**hence**  
 $\vdash \sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) \Psi) \leftrightarrow (?\chi' \sqcap \sqcup (\text{map } (\lambda\varphi. ?\Pi\Delta \varphi) \Psi))$   
**using** *conj-extract by presburger*  
**moreover have**  
 $\vdash \sqcup (\text{map } ?\Pi\Delta \Psi) \leftrightarrow (\sqcap ?\Delta' \sqcap \sqcup (\text{map } ?\Pi \Psi))$   
 $\rightarrow \sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) \Psi) \leftrightarrow (?\chi' \sqcap \sqcup (\text{map } ?\Pi\Delta \Psi))$   
 $\rightarrow \sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) \Psi) \leftrightarrow ((?\chi' \sqcap \sqcap ?\Delta') \sqcap \sqcup (\text{map } ?\Pi \Psi))$   
**proof** –  
**let**  $? \varphi = \langle \sqcup (\text{map } ?\Pi\Delta \Psi) \rangle \leftrightarrow (\langle \sqcap ?\Delta' \rangle \sqcap \langle \sqcup (\text{map } ?\Pi \Psi) \rangle)$   
 $\rightarrow \langle \sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) \Psi) \rangle \leftrightarrow (\langle ?\chi' \rangle \sqcap \langle \sqcup (\text{map } ?\Pi\Delta \Psi) \rangle)$   
 $\rightarrow \langle \sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap ?\Pi\Delta \varphi) \Psi) \rangle \leftrightarrow ((\langle ?\chi' \rangle \sqcap \langle \sqcap ?\Delta' \rangle) \sqcap \langle \sqcup (\text{map } ?\Pi \Psi) \rangle)$   
*(map ?Pi Psi))*  
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  **by** *fastforce*  
**hence**  $\vdash \langle ?\varphi \rangle$  **using** *propositional-semantic by blast*  
**thus** *?thesis by simp*  
**qed**  
**ultimately have**  $\vdash \sqcup (\text{map } (\lambda\varphi. ?\chi' \sqcap \sqcap (?\Delta' @ \text{map } \text{snd } \varphi)) \Psi)$   
 $\leftrightarrow ((?\chi' \sqcap \sqcap ?\Delta') \sqcap \sqcup (\text{map } (\lambda\varphi. \sqcap (\text{map } \text{snd } \varphi)) \Psi))$   
**using** *modus-ponens by blast*  
**thus** *?case by simp*  
**qed**

**lemma** (*in classical-logic*) *extract-inner-concat-remdups*:  
 $\vdash \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups} \circ (@) \Delta)) \Psi) \leftrightarrow$   
 $(\sqcap (\text{map } \text{snd } \Delta) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups})) \Psi))$   
**proof** –  
**have**  $\forall \Psi. \vdash \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups} \circ (@) \Delta)) \Psi) \leftrightarrow$   
 $(\sqcap (\text{map } \text{snd } \Delta) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups})) \Psi))$   
**proof** (*induct*  $\Delta$ )  
**case** *Nil*  
**then show** *?case*  
**by** (*simp,*  
*meson modus-ponens*  
*biconditional-introduction*  
*conjunction-introduction*  
*conjunction-right-elimination*  
*verum-tautology*)

**next**  
**case** (*Cons*  $\delta \Delta$ )  
 $\{$

```

fix  $\Psi$ 
have  $\vdash \sqcup (map (\sqcap \circ (map\ snd \circ remdups \circ (@) (\delta \# \Delta))) \Psi)$ 
       $\leftrightarrow (\sqcap (map\ snd (\delta \# \Delta)) \sqcap \sqcup (map (\sqcap \circ (map\ snd \circ remdups)) \Psi))$ 
proof (cases  $\delta \in set\ \Delta$ )
  assume  $\delta \in set\ \Delta$ 
  have
     $\vdash \sqcap (map\ snd\ \Delta) \leftrightarrow (snd\ \delta \sqcap \sqcap (map\ snd\ \Delta))$ 
     $\rightarrow \sqcup (map (\sqcap \circ (map\ snd \circ remdups \circ (@) \Delta)) \Psi)$ 
     $\leftrightarrow (\sqcap (map\ snd\ \Delta) \sqcap \sqcup (map (\sqcap \circ (map\ snd \circ remdups)) \Psi))$ 
     $\rightarrow \sqcup (map (\sqcap \circ (map\ snd \circ remdups \circ (@) \Delta)) \Psi)$ 
     $\leftrightarrow ((snd\ \delta \sqcap \sqcap (map\ snd\ \Delta)) \sqcap \sqcup (map (\sqcap \circ (map\ snd \circ remdups))$ 
 $\Psi))$ 
  proof -
    let  $?\varphi = \langle \sqcap (map\ snd\ \Delta) \rangle \leftrightarrow (\langle snd\ \delta \rangle \sqcap \langle \sqcap (map\ snd\ \Delta) \rangle)$ 
       $\rightarrow \langle \sqcup (map (\sqcap \circ (map\ snd \circ remdups \circ (@) \Delta)) \Psi) \rangle$ 
       $\leftrightarrow (\langle \sqcap (map\ snd\ \Delta) \rangle \sqcap \langle \sqcup (map (\sqcap \circ (map\ snd \circ remdups))$ 
 $\Psi) \rangle)$ 
       $\rightarrow \langle \sqcup (map (\sqcap \circ (map\ snd \circ remdups \circ (@) \Delta)) \Psi) \rangle$ 
       $\leftrightarrow ((\langle snd\ \delta \rangle \sqcap \langle \sqcap (map\ snd\ \Delta) \rangle) \sqcap \langle \sqcup (map (\sqcap \circ (map\ snd \circ$ 
 $remdups)) \Psi) \rangle)$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  by fastforce
    hence  $\vdash (\ ?\varphi \ )$  using propositional-semantic by blast
    thus  $?thesis$  by simp
  qed
moreover have  $\vdash \sqcap (map\ snd\ \Delta) \leftrightarrow (snd\ \delta \sqcap \sqcap (map\ snd\ \Delta))$ 
  by (simp add:  $\langle \delta \in set\ \Delta \rangle$  conj-absorption)
ultimately have
   $\vdash \sqcup (map (\sqcap \circ (map\ snd \circ remdups \circ (@) \Delta)) \Psi)$ 
     $\leftrightarrow ((snd\ \delta \sqcap \sqcap (map\ snd\ \Delta)) \sqcap \sqcup (map (\sqcap \circ (map\ snd \circ remdups))$ 
 $\Psi))$ 
  using Cons.hyps modus-ponens by blast
moreover have  $map\ snd \circ remdups \circ (@) (\delta \# \Delta) = map\ snd \circ remdups$ 
 $\circ (@) \Delta$ 
  using  $\langle \delta \in set\ \Delta \rangle$  by fastforce
ultimately show  $?thesis$  using Cons by simp
next
assume  $\delta \notin set\ \Delta$ 
hence  $\dagger$ :
   $\sqcap \circ (map\ snd \circ remdups) = (\lambda\psi. \sqcap (map\ snd (remdups\ \psi)))$ 
   $(\lambda\psi. \sqcap (map\ snd (if\ \delta \in set\ \psi\ then\ remdups\ (\Delta\ @\ \psi)\ else\ \delta \# remdups$ 
 $(\Delta\ @\ \psi))))$ 
   $= \sqcap \circ (map\ snd \circ remdups \circ (@) (\delta \# \Delta))$ 
  by fastforce+
show  $?thesis$ 
proof (induct  $\Psi$ )
  case Nil
  then show  $?case$ 
  by (simp, metis list.simps(8) arbitrary-disjunction.simps(1) conj-extract)
next

```



```

case (Cons  $\psi$   $\Psi$ )
have  $\vdash \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) \Delta)) [\psi])$ 
       $\leftrightarrow (\sqcap (\text{map snd } \Delta) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups})) [\psi]))$ 
using  $\langle \forall \Psi. \vdash \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) \Delta)) \Psi)$ 
       $\leftrightarrow (\sqcap (\text{map snd } \Delta) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups}))$ 
 $\Psi)) \rangle$ 
by blast
hence
 $\vdash (\sqcap (\text{map snd } (\text{remdups } (\Delta @ \psi))) \sqcup \perp)$ 
 $\leftrightarrow (\sqcap (\text{map snd } \Delta) \sqcap \sqcap (\text{map snd } (\text{remdups } \psi)) \sqcup \perp)$ 
by simp
hence  $\star$ :
 $\vdash \sqcap (\text{map snd } (\text{remdups } (\Delta @ \psi))) \leftrightarrow (\sqcap (\text{map snd } \Delta) \sqcap \sqcap (\text{map snd } (\text{remdups } \psi)))$ 
by (metis
      (no-types, opaque-lifting)
      (biconditional-conjunction-weaken-rule)
      (biconditional-symmetry-rule)
      (biconditional-transitivity-rule)
      (disjunction-def)
      (double-negation-biconditional)
      (negation-def))
have  $\vdash \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)$ 
       $\leftrightarrow (\sqcap (\text{map snd } (\delta \# \Delta)) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups}))$ 
 $\Psi))$ 
using Cons by blast
hence  $\diamond$ :  $\vdash \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)$ 
       $\leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd } \Delta)) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$ 
 $\text{remdups})) \Psi))$ 
by simp
show ?case
proof (cases  $\delta \in \text{set } \psi$ )
assume  $\delta \in \text{set } \psi$ 
have  $\text{snd } \delta \in \text{set } (\text{map snd } (\text{remdups } \psi))$ 
using  $\langle \delta \in \text{set } \psi \rangle$  by auto
hence  $\spadesuit$ :  $\vdash \sqcap (\text{map snd } (\text{remdups } \psi)) \leftrightarrow (\text{snd } \delta \sqcap \sqcap (\text{map snd } (\text{remdups } \psi)))$ 
using conj-absorption by blast
have
 $\vdash (\sqcap (\text{map snd } (\text{remdups } \psi)) \leftrightarrow (\text{snd } \delta \sqcap \sqcap (\text{map snd } (\text{remdups } \psi))))$ 
 $\rightarrow (\sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)$ 
 $\leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd } \Delta)) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$ 
 $\text{remdups})) \Psi)))$ 
 $\rightarrow (\sqcap (\text{map snd } (\text{remdups } (\Delta @ \psi))) \leftrightarrow (\sqcap (\text{map snd } \Delta) \sqcap \sqcap (\text{map$ 
 $\text{snd } (\text{remdups } \psi))))$ 
 $\rightarrow (\sqcap (\text{map snd } (\text{remdups } (\Delta @ \psi)))$ 
 $\sqcup \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi))$ 
 $\leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd } \Delta))$ 

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$\sqcap (\sqcap (\text{map snd (remdups } \psi)) \sqcup \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi)))$   
**proof** –  
**let**  $? \varphi =$   
 $(\langle \sqcap (\text{map snd (remdups } \psi) \rangle \leftrightarrow (\langle \text{snd } \delta \rangle \sqcap \langle \sqcap (\text{map snd (remdups}$   
 $\psi) \rangle)))$   
 $\rightarrow (\langle \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi) \rangle$   
 $\leftrightarrow ((\langle \text{snd } \delta \rangle \sqcap \langle \sqcap (\text{map snd } \Delta) \rangle) \sqcap \langle \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi) \rangle))$   
 $\rightarrow (\langle \sqcap (\text{map snd (remdups } (\Delta @ \psi)) \rangle$   
 $\leftrightarrow (\langle \sqcap (\text{map snd } \Delta) \rangle \sqcap \langle \sqcap (\text{map snd (remdups } \psi) \rangle))$   
 $\rightarrow (\langle \sqcap (\text{map snd (remdups } (\Delta @ \psi)) \rangle$   
 $\sqcup \langle \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi) \rangle)$   
 $\leftrightarrow ((\langle \text{snd } \delta \rangle \sqcap \langle \sqcap (\text{map snd } \Delta) \rangle)$   
 $\sqcap (\langle \sqcap (\text{map snd (remdups } \psi) \rangle \sqcup \langle \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi) \rangle))$   
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  **by fastforce**  
**hence**  $\vdash (\langle ? \varphi \rangle)$  **using propositional-semantic by blast**  
**thus**  $?thesis$  **by simp**  
**qed**  
**hence**  
 $\vdash (\sqcap (\text{map snd (remdups } (\Delta @ \psi))$   
 $\sqcup \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)$   
 $\leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd } \Delta))$   
 $\sqcap (\sqcap (\text{map snd (remdups } \psi)) \sqcup \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi)))$   
**using**  $\star \diamond \spadesuit$  **modus-ponens by blast**  
**thus**  $?thesis$  **using**  $\langle \delta \notin \text{set } \Delta \rangle \langle \delta \in \text{set } \psi \rangle$   
**by (simp add: †)**  
**next**  
**assume**  $\delta \notin \text{set } \psi$   
**have**  
 $\vdash (\sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)$   
 $\leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd } \Delta)) \sqcap \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi)))$   
 $\rightarrow (\sqcap (\text{map snd (remdups } (\Delta @ \psi)) \leftrightarrow (\sqcap (\text{map snd } \Delta) \sqcap \sqcap (\text{map}$   
 $\text{snd (remdups } \psi)))$   
 $\rightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd (remdups } (\Delta @ \psi)))$   
 $\sqcup \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)$   
 $\leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map snd } \Delta))$   
 $\sqcap (\sqcap (\text{map snd (remdups } \psi)) \sqcup \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi)))$   
**proof** –  
**let**  $? \varphi =$   
 $(\langle \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi) \rangle$   
 $\leftrightarrow ((\langle \text{snd } \delta \rangle \sqcap \langle \sqcap (\text{map snd } \Delta) \rangle) \sqcap \langle \sqcup (\text{map } (\sqcap \circ (\text{map snd} \circ$   
 $\text{remdups})) \Psi) \rangle))$   
 $\rightarrow (\langle \sqcap (\text{map snd (remdups } (\Delta @ \psi)) \rangle$   
 $\leftrightarrow (\langle \sqcap (\text{map snd } \Delta) \rangle \sqcap \langle \sqcap (\text{map snd (remdups } \psi) \rangle))$

$\rightarrow ((\langle \text{snd } \delta \rangle \sqcap \langle \sqcap (\text{map } \text{snd } (\text{remdups } (\Delta @ \psi))) \rangle) \sqcup \langle \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi) \rangle)$   
 $\leftrightarrow ((\langle \text{snd } \delta \rangle \sqcap \langle \sqcap (\text{map } \text{snd } \Delta) \rangle) \sqcap \langle \sqcap (\text{map } \text{snd } (\text{remdups } \psi)) \rangle \sqcup \langle \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups})) \Psi) \rangle)$   
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  **by** *fastforce*  
**hence**  $\vdash (\langle ?\varphi \rangle)$  **using** *propositional-semantic* **by** *blast*  
**thus** *?thesis* **by** *simp*  
**qed**  
**hence**  
 $\vdash ((\text{snd } \delta \sqcap \sqcap (\text{map } \text{snd } (\text{remdups } (\Delta @ \psi))) \sqcup \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups} \circ (@) (\delta \# \Delta))) \Psi)) \leftrightarrow ((\text{snd } \delta \sqcap \sqcap (\text{map } \text{snd } \Delta)) \sqcap (\sqcap (\text{map } \text{snd } (\text{remdups } \psi)) \sqcup \sqcup (\text{map } (\sqcap \circ (\text{map } \text{snd} \circ \text{remdups})) \Psi)))$   
**using**  $\star \diamond$  *modus-ponens* **by** *blast*  
**then show** *?thesis* **using**  $\langle \delta \notin \text{set } \psi \rangle \langle \delta \notin \text{set } \Delta \rangle$  **by** (*simp add: †*)  
**qed**  
**qed**  
**qed**  
**}**  
**then show** *?case* **by** *fastforce*  
**qed**  
**thus** *?thesis* **by** *blast*  
**qed**

**lemma** (in *classical-logic*) *optimal-witness-list-intersect-biconditional*:

**assumes**  $\text{mset } \Xi \subseteq\# \text{mset } \Gamma$   
**and**  $\text{mset } \Phi \subseteq\# \text{mset } (\Gamma \ominus \Xi)$   
**and**  $\text{mset } \Psi \subseteq\# \text{mset } (\mathfrak{M} \rightarrow \varphi \Xi)$   
**shows**  $\exists \Sigma. \vdash ((\Phi @ \Psi) \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \Sigma) \rightarrow \varphi) \wedge (\forall \sigma \in \text{set } \Sigma. \text{mset } \sigma \subseteq\# \text{mset } \Gamma \wedge \text{length } \sigma + 1 \geq \text{length } (\Phi @ \Psi))$   
**proof** –  
**have**  $\exists \Sigma. \vdash (\Psi \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \Sigma) \rightarrow \varphi) \wedge (\forall \sigma \in \text{set } \Sigma. \text{mset } \sigma \subseteq\# \text{mset } \Xi \wedge \text{length } \sigma + 1 \geq \text{length } \Psi)$   
**proof** –  
**from** *assms*( $\beta$ ) **obtain**  $\Psi_0 :: ('a \text{ list } \times 'a) \text{ list}$  **where**  $\Psi_0$ :  
 $\text{mset } \Psi_0 \subseteq\# \text{mset } (\mathfrak{M} \Xi)$   
 $\text{map } (\lambda(\Psi, \psi). (\Psi \rightarrow \varphi \rightarrow \psi)) \Psi_0 = \Psi$   
**using** *mset-sub-map-list-exists* **by** *fastforce*  
**let**  $? \Pi_C = \lambda (\Delta, \delta) \Sigma. (\text{map } ((\#) (\Delta, \delta)) \Sigma) @ (\text{map } ((@) (\mathfrak{M} \Delta)) \Sigma)$   
**let**  $?T_\Sigma = \lambda \Psi. \text{foldr } ? \Pi_C \Psi []$   
**let**  $? \Sigma = \text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma \Psi_0)$   
**have**  $I: \vdash (\Psi \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap ? \Sigma) \rightarrow \varphi)$   
**proof** –  
**let**  $? \Sigma_\alpha = \text{map } (\text{map } \text{snd}) (?T_\Sigma \Psi_0)$   
**let**  $? \Psi' = \text{map } (\lambda(\Psi, \psi). (\Psi \rightarrow \varphi \rightarrow \psi)) \Psi_0$   
**{**

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fix  $\Psi :: ('a \text{ list} \times 'a) \text{ list}$ 
let  $? \Sigma_\alpha = \text{map } (\text{map } \text{snd}) (?T_\Sigma \Psi)$ 
let  $? \Sigma = \text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma \Psi)$ 
have  $\vdash (\bigsqcup (\text{map } \sqcap ? \Sigma_\alpha) \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap ? \Sigma) \rightarrow \varphi)$ 
proof (induct  $\Psi$ )
  case Nil
  then show ?case by (simp add: biconditional-reflection)
next
  case (Cons  $\Delta \delta \Psi$ )
  let  $? \Delta = \text{fst } \Delta \delta$ 
  let  $? \delta = \text{snd } \Delta \delta$ 
  let  $? \Sigma_\alpha = \text{map } (\text{map } \text{snd}) (?T_\Sigma \Psi)$ 
  let  $? \Sigma = \text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma \Psi)$ 
  let  $? \Sigma'_\alpha = \text{map } (\text{map } \text{snd}) (?T_\Sigma ((? \Delta, ? \delta) \# \Psi))$ 
  let  $? \Sigma' = \text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma ((? \Delta, ? \delta) \# \Psi))$ 
  {
    fix  $\Delta :: 'a \text{ list}$ 
    fix  $\delta :: 'a$ 
    let  $? \Sigma'_\alpha = \text{map } (\text{map } \text{snd}) (?T_\Sigma ((\Delta, \delta) \# \Psi))$ 
    let  $? \Sigma' = \text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma ((\Delta, \delta) \# \Psi))$ 
    let  $? \Phi = \text{map } (\text{map } \text{snd} \circ (@) [(\Delta, \delta)]) (?T_\Sigma \Psi)$ 
    let  $? \Psi = \text{map } (\text{map } \text{snd} \circ (@) (\mathfrak{V} \Delta)) (?T_\Sigma \Psi)$ 
    let  $? \Delta = \text{map } (\text{map } \text{snd} \circ \text{remdups} \circ (@) [(\Delta, \delta)]) (?T_\Sigma \Psi)$ 
    let  $? \Omega = \text{map } (\text{map } \text{snd} \circ \text{remdups} \circ (@) (\mathfrak{V} \Delta)) (?T_\Sigma \Psi)$ 
    have  $\vdash (\bigsqcup (\text{map } \sqcap ? \Phi @ \text{map } \sqcap ? \Psi) \leftrightarrow (\bigsqcup (\text{map } \sqcap ? \Phi) \sqcup \bigsqcup (\text{map } \sqcap ? \Psi))) \rightarrow$ 
     $(\bigsqcup (\text{map } \sqcap ? \Delta @ \text{map } \sqcap ? \Omega) \leftrightarrow (\bigsqcup (\text{map } \sqcap ? \Delta) \sqcup \bigsqcup (\text{map } \sqcap ? \Omega))) \rightarrow$ 
     $(\bigsqcup (\text{map } \sqcap ? \Phi) \leftrightarrow (\sqcap [\delta] \sqcap \bigsqcup (\text{map } \sqcap ? \Sigma_\alpha))) \rightarrow$ 
     $(\bigsqcup (\text{map } \sqcap ? \Psi) \leftrightarrow (\sqcap \Delta \sqcap \bigsqcup (\text{map } \sqcap ? \Sigma_\alpha))) \rightarrow$ 
     $(\bigsqcup (\text{map } \sqcap ? \Delta) \leftrightarrow (\sqcap [\delta] \sqcap \bigsqcup (\text{map } \sqcap ? \Sigma))) \rightarrow$ 
     $(\bigsqcup (\text{map } \sqcap ? \Omega) \leftrightarrow (\sqcap \Delta \sqcap \bigsqcup (\text{map } \sqcap ? \Sigma))) \rightarrow$ 
     $((\bigsqcup (\text{map } \sqcap ? \Sigma_\alpha) \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap ? \Sigma) \rightarrow \varphi)) \rightarrow$ 
     $((\bigsqcup (\text{map } \sqcap ? \Phi @ \text{map } \sqcap ? \Psi) \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap ? \Delta @ \text{map } \sqcap ? \Omega) \rightarrow \varphi))$ 
    proof –
    let  $? \varphi =$ 
     $(\langle \bigsqcup (\text{map } \sqcap ? \Phi @ \text{map } \sqcap ? \Psi) \rangle \leftrightarrow (\langle \bigsqcup (\text{map } \sqcap ? \Phi) \rangle \sqcup \langle \bigsqcup (\text{map } \sqcap ? \Psi) \rangle)) \rightarrow$ 
     $(\langle \bigsqcup (\text{map } \sqcap ? \Delta @ \text{map } \sqcap ? \Omega) \rangle \leftrightarrow (\langle \bigsqcup (\text{map } \sqcap ? \Delta) \rangle \sqcup \langle \bigsqcup (\text{map } \sqcap ? \Omega) \rangle)) \rightarrow$ 
     $(\langle \bigsqcup (\text{map } \sqcap ? \Phi) \rangle \leftrightarrow (\langle \sqcap [\delta] \rangle \sqcap \langle \bigsqcup (\text{map } \sqcap ? \Sigma_\alpha) \rangle)) \rightarrow$ 
     $(\langle \bigsqcup (\text{map } \sqcap ? \Psi) \rangle \leftrightarrow (\langle \sqcap \Delta \rangle \sqcap \langle \bigsqcup (\text{map } \sqcap ? \Sigma_\alpha) \rangle)) \rightarrow$ 
     $(\langle \bigsqcup (\text{map } \sqcap ? \Delta) \rangle \leftrightarrow (\langle \sqcap [\delta] \rangle \sqcap \langle \bigsqcup (\text{map } \sqcap ? \Sigma) \rangle)) \rightarrow$ 
     $(\langle \bigsqcup (\text{map } \sqcap ? \Omega) \rangle \leftrightarrow (\langle \sqcap \Delta \rangle \sqcap \langle \bigsqcup (\text{map } \sqcap ? \Sigma) \rangle)) \rightarrow$ 
     $((\langle \bigsqcup (\text{map } \sqcap ? \Sigma_\alpha) \rangle \rightarrow \langle \varphi \rangle) \leftrightarrow (\langle \bigsqcup (\text{map } \sqcap ? \Sigma) \rangle \rightarrow \langle \varphi \rangle)) \rightarrow$ 
     $((\langle \bigsqcup (\text{map } \sqcap ? \Phi @ \text{map } \sqcap ? \Psi) \rangle \rightarrow \langle \varphi \rangle) \leftrightarrow (\langle \bigsqcup (\text{map } \sqcap ? \Delta @ \text{map } \sqcap ? \Omega) \rangle \rightarrow \langle \varphi \rangle))$ 
    have  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  by fastforce
  }

```

hence  $\vdash (\ ?\varphi \ )$  **using** *propositional-semantic* **by** *blast*  
 thus *?thesis* **by** *simp*  
**qed**  
**moreover**  
**have** *map snd*  $(\mathfrak{V} \Delta) = \Delta$  **by** (*induct*  $\Delta$ , *auto*)  
**hence**  $\vdash \sqcup (\text{map } \sqcap \ ?\Phi @ \text{map } \sqcap \ ?\Psi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Phi) \sqcup \sqcup (\text{map } \sqcap \ ?\Psi))$   
 $\vdash \sqcup (\text{map } \sqcap \ ?\Delta @ \text{map } \sqcap \ ?\Omega) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Delta) \sqcup \sqcup (\text{map } \sqcap \ ?\Omega))$   
 $\vdash \sqcup (\text{map } \sqcap \ ?\Phi) \leftrightarrow (\sqcap [\delta] \sqcap \sqcup (\text{map } \sqcap \ ?\Sigma_\alpha))$   
 $\vdash \sqcup (\text{map } \sqcap \ ?\Psi) \leftrightarrow (\sqcap \Delta \sqcap \sqcup (\text{map } \sqcap \ ?\Sigma_\alpha))$   
 $\vdash \sqcup (\text{map } \sqcap \ ?\Delta) \leftrightarrow (\sqcap [\delta] \sqcap \sqcup (\text{map } \sqcap \ ?\Sigma))$   
 $\vdash \sqcup (\text{map } \sqcap \ ?\Omega) \leftrightarrow (\sqcap \Delta \sqcap \sqcup (\text{map } \sqcap \ ?\Sigma))$   
**using** *arbitrary-disj-concat-equiv*  
*extract-inner-concat* [**where**  $\Delta = [(\Delta, \delta)]$  **and**  $\Psi = ?T_\Sigma \Psi$ ]  
*extract-inner-concat* [**where**  $\Delta = \mathfrak{V} \Delta$  **and**  $\Psi = ?T_\Sigma \Psi$ ]  
*extract-inner-concat-remdups* [**where**  $\Delta = [(\Delta, \delta)]$  **and**  $\Psi = ?T_\Sigma \Psi$ ]  
 $\vdash \sqcup (\text{map } \sqcap \ ?\Delta @ \text{map } \sqcap \ ?\Omega) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Delta) \sqcup \sqcup (\text{map } \sqcap \ ?\Omega))$   
*extract-inner-concat-remdups* [**where**  $\Delta = \mathfrak{V} \Delta$  **and**  $\Psi = ?T_\Sigma \Psi$ ]  
**by** *auto*  
**ultimately have**  
 $\vdash ((\sqcup (\text{map } \sqcap \ ?\Sigma_\alpha) \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Sigma) \rightarrow \varphi)) \rightarrow$   
 $(\sqcup (\text{map } \sqcap \ ?\Phi @ \text{map } \sqcap \ ?\Psi) \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Delta @ \text{map } \sqcap \ ?\Omega) \rightarrow \varphi)$   
**using** *modus-ponens* **by** *blast*  
**moreover have**  $(\#) (\Delta, \delta) = (@) [(\Delta, \delta)]$  **by** *fastforce*  
**ultimately have**  
 $\vdash ((\sqcup (\text{map } \sqcap \ ?\Sigma_\alpha) \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Sigma) \rightarrow \varphi)) \rightarrow$   
 $((\sqcup (\text{map } \sqcap \ ?\Sigma_\alpha') \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Sigma') \rightarrow \varphi))$   
**by** *auto*  
**}**  
**hence**  
 $\vdash ((\sqcup (\text{map } \sqcap \ ?\Sigma_\alpha') \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Sigma') \rightarrow \varphi))$   
**using** *Cons modus-ponens* **by** *blast*  
**moreover have**  $\Delta\delta = (? \Delta, ? \delta)$  **by** *fastforce*  
**ultimately show** *?case* **by** *metis*  
**qed**  
**}**  
**hence**  $\vdash (\sqcup (\text{map } \sqcap \ ?\Sigma_\alpha) \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Sigma) \rightarrow \varphi)$  **by** *blast*  
**moreover have**  $\vdash (? \Psi' : \rightarrow \varphi) \leftrightarrow (\sqcup (\text{map } \sqcap \ ?\Sigma_\alpha) \rightarrow \varphi)$   
**proof** (*induct*  $\Psi_0$ )  
**case** *Nil*  
**have**  $\vdash \varphi \leftrightarrow ((\top \sqcup \perp) \rightarrow \varphi)$   
**proof** –  
**let**  $? \varphi = \langle \varphi \rangle \leftrightarrow ((\top \sqcup \perp) \rightarrow \langle \varphi \rangle)$   
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ? \varphi$  **by** *fastforce*  
**hence**  $\vdash (\ ?\varphi \ )$  **using** *propositional-semantic* **by** *blast*  
**thus** *?thesis* **by** *simp*  
**qed**

```

thus ?case by simp
next
case (Cons  $\psi_0$   $\Psi_0$ )
let ? $\Xi$  = fst  $\psi_0$ 
let ? $\delta$  = snd  $\psi_0$ 
let ? $\Psi'$  = map ( $\lambda(\Psi, \psi). (\Psi \rightarrow \varphi \rightarrow \psi)$ )  $\Psi_0$ 
let ? $\Sigma_\alpha$  = map (map snd) (? $T_\Sigma$   $\Psi_0$ )
{
  fix  $\Xi :: 'a$  list
  have map snd ( $\mathfrak{V} \Xi$ ) =  $\Xi$  by (induct  $\Xi$ , auto)
  hence map snd  $\circ$  (@) ( $\mathfrak{V} \Xi$ ) = (@)  $\Xi$   $\circ$  map snd by fastforce
}
moreover have (map snd  $\circ$  (#) (? $\Xi$ , ? $\delta$ )) = (@) [? $\delta$ ]  $\circ$  map snd by fastforce
ultimately have †:
  map (map snd) (? $T_\Sigma$  ( $\psi_0$  #  $\Psi_0$ )) = map ((#) ? $\delta$ ) ? $\Sigma_\alpha$  @ map ((@) ? $\Xi$ )
? $\Sigma_\alpha$ 
  map ( $\lambda(\Psi, \psi). (\Psi \rightarrow \varphi \rightarrow \psi)$ ) ( $\psi_0$  #  $\Psi_0$ ) = ? $\Xi$   $\rightarrow$   $\varphi$   $\rightarrow$  ? $\delta$  # ? $\Psi'$ 
  by (simp add: case-prod-beta')+
have A: † (? $\Psi' \rightarrow \varphi$ )  $\leftrightarrow$  ( $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ )  $\rightarrow$   $\varphi$ ) using Cons.hyps by auto
have B: † (? $\Xi \rightarrow \varphi$ )  $\leftrightarrow$  ( $\prod$  ? $\Xi \rightarrow \varphi$ )
  by (simp add: list-curry-uncurry)
have C: †  $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ) @ map  $\sqcap$  (map ((@) ? $\Xi$ )
? $\Sigma_\alpha$ ))
   $\leftrightarrow$  ( $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ))  $\sqcup$   $\bigsqcup$  (map  $\sqcap$  (map ((@) ? $\Xi$ )
? $\Sigma_\alpha$ )))
  using arbitrary-disj-concat-equiv by blast
have map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ) = (map (( $\prod$ ) ? $\delta$ ) (map  $\sqcap$  ? $\Sigma_\alpha$ )) by auto
hence D: †  $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ))  $\leftrightarrow$  (? $\delta$   $\sqcap$   $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ ))
  using conj-extract by presburger
have E: †  $\bigsqcup$  (map  $\sqcap$  (map ((@) ? $\Xi$ ) ? $\Sigma_\alpha$ ))  $\leftrightarrow$  ( $\prod$  ? $\Xi$   $\sqcap$   $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ ))
  using conj-multi-extract by blast
have
  †
  (? $\Psi' \rightarrow \varphi$ )  $\leftrightarrow$  ( $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ )  $\rightarrow$   $\varphi$ )
   $\rightarrow$  (? $\Xi \rightarrow \varphi$ )  $\leftrightarrow$  ( $\prod$  ? $\Xi \rightarrow \varphi$ )
   $\rightarrow$   $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ) @ map  $\sqcap$  (map ((@) ? $\Xi$ ) ? $\Sigma_\alpha$ )
   $\leftrightarrow$  ( $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ))  $\sqcup$   $\bigsqcup$  (map  $\sqcap$  (map ((@) ? $\Xi$ )
? $\Sigma_\alpha$ )))
   $\rightarrow$   $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ))  $\leftrightarrow$  (? $\delta$   $\sqcap$   $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ ))
   $\rightarrow$   $\bigsqcup$  (map  $\sqcap$  (map ((@) ? $\Xi$ ) ? $\Sigma_\alpha$ ))  $\leftrightarrow$  ( $\prod$  ? $\Xi$   $\sqcap$   $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ ))
   $\rightarrow$  ((? $\Xi \rightarrow \varphi \rightarrow ?\delta$ )  $\rightarrow$  ? $\Psi' \rightarrow \varphi$ )
   $\leftrightarrow$  ( $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ) @ map  $\sqcap$  (map ((@) ? $\Xi$ ) ? $\Sigma_\alpha$ ))
   $\rightarrow$   $\varphi$ )
proof –
  let ? $\varphi$  =
   $\langle$  ? $\Psi' \rightarrow \varphi$   $\rangle$   $\leftrightarrow$  ( $\langle$   $\bigsqcup$  (map  $\sqcap$  ? $\Sigma_\alpha$ )  $\rangle$   $\rightarrow$   $\langle$   $\varphi$   $\rangle$ )
   $\rightarrow$   $\langle$  (? $\Xi \rightarrow \varphi$ )  $\rangle$   $\leftrightarrow$  ( $\langle$   $\prod$  ? $\Xi$   $\rangle$   $\rightarrow$   $\langle$   $\varphi$   $\rangle$ )
   $\rightarrow$   $\langle$   $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ) @ map  $\sqcap$  (map ((@) ? $\Xi$ )
? $\Sigma_\alpha$ ))  $\rangle$ 
   $\leftrightarrow$  ( $\langle$   $\bigsqcup$  (map  $\sqcap$  (map ((#) ? $\delta$ ) ? $\Sigma_\alpha$ ))  $\rangle$   $\sqcup$   $\langle$   $\bigsqcup$  (map  $\sqcap$  (map ((@)

```

```

(⟨∃⟩ ?Σα)))
  → ⟨⊔ (map ⊓ (map ((#) ?δ) ?Σα))⟩ ↔ (⟨?δ⟩ ⊓ ⟨⊔ (map ⊓
?Σα))⟩
  → ⟨⊔ (map ⊓ (map ((@) ?∃) ?Σα))⟩ ↔ (⟨⊓ ?∃⟩ ⊓ ⟨⊔ (map ⊓
?Σα))⟩
  → ((⟨?∃ :→ φ⟩ → ⟨?δ⟩) → ⟨?Ψ' :→ φ⟩)
    ↔ (⟨⊔ (map ⊓ (map ((#) ?δ) ?Σα) @ map ⊓ (map ((@) ?∃)
?Σα))⟩ → ⟨φ⟩)
  have ∀ M. M ⊨prop ?φ by fastforce
  hence ⊢ ( ?φ ) using propositional-semantic by blast
  thus ?thesis by simp
qed
hence
  ⊢ ((?∃ :→ φ → ?δ) → ?Ψ' :→ φ)
    ↔ (⊔ (map ⊓ (map ((#) ?δ) ?Σα) @ map ⊓ (map ((@) ?∃) ?Σα)) →
φ)
  using A B C D E modus-ponens by blast
  thus ?case using † by simp
qed
ultimately show ?thesis using biconditional-transitivity-rule Ψ0 by blast
qed
have II: ∀ σ ∈ set ?Σ. length σ + 1 ≥ length Ψ
proof -
  let ?F = length ∘ fst
  let ?S = sort-key (- ?F)
  let ?Σ' = map (map snd ∘ remdups) (?TΣ (?S Ψ0))
  have mset Ψ0 = mset (?S Ψ0) by simp

  have ∀ Φ. mset Ψ0 = mset Φ → mset (map mset (?TΣ Ψ0)) = mset (map
mset (?TΣ Φ))
  proof (induct Ψ0)
  case Nil
  then show ?case by simp
  next
  case (Cons ψ Ψ0)
  obtain Δ δ where ψ = (Δ, δ) by fastforce
  {
  fix Φ
  assume mset (ψ # Ψ0) = mset Φ
  hence mset Ψ0 = mset (remove1 ψ Φ)
  by (simp add: union-single-eq-diff)
  have ψ ∈ set Φ using ⟨mset (ψ # Ψ0) = mset Φ⟩
  by (metis list.set-intros(1) set-mset-mset)
  hence mset (map mset (?TΣ Φ)) = mset (map mset (?TΣ (ψ # (remove1
ψ Φ))))
  proof (induct Φ)
  case Nil
  then show ?case by simp
  next

```

```

case (Cons  $\varphi$   $\Phi$ )
then show ?case proof (cases  $\varphi = \psi$ )
  case True
    then show ?thesis by simp
  next
    case False
      let ? $\Sigma'$  = ?T $\Sigma$  ( $\psi$  # (remove1  $\psi$   $\Phi$ ))
      have †: mset (map mset ? $\Sigma'$ ) = mset (map mset (?T $\Sigma$   $\Phi$ ))
        using Cons False by simp
      obtain  $\Delta'$   $\delta'$ 
        where  $\varphi = (\Delta', \delta')$ 
        by fastforce
      let ? $\Sigma$  = ?T $\Sigma$  (remove1  $\psi$   $\Phi$ )
      let ?m = image-mset mset
      have
        mset (map mset (?T $\Sigma$  ( $\psi$  # remove1  $\psi$  ( $\varphi$  #  $\Phi$ )))) =
          mset (map mset (? $\Pi_C$   $\psi$  (? $\Pi_C$   $\varphi$  ? $\Sigma$ )))
        using False by simp
      hence mset (map mset (?T $\Sigma$  ( $\psi$  # remove1  $\psi$  ( $\varphi$  #  $\Phi$ )))) =
        (?m  $\circ$  image-mset ((#)  $\psi$ )  $\circ$  image-mset ((#)  $\varphi$ )) (mset ? $\Sigma$ ) +
        (?m  $\circ$  image-mset ((#)  $\psi$ )  $\circ$  image-mset (@) ( $\mathfrak{V}$   $\Delta'$ ))) (mset
?Σ) +
        (?m  $\circ$  (image-mset (@) ( $\mathfrak{V}$   $\Delta$ ))  $\circ$  image-mset ((#)  $\varphi$ )) (mset
?Σ) +
        (?m  $\circ$  (image-mset (@) ( $\mathfrak{V}$   $\Delta$ ))  $\circ$  image-mset (@) ( $\mathfrak{V}$   $\Delta'$ )))
(mset ?Σ)
        using  $\langle \psi = (\Delta, \delta) \rangle$   $\langle \varphi = (\Delta', \delta') \rangle$ 
        by (simp add: multiset.map-comp)
      hence mset (map mset (?T $\Sigma$  ( $\psi$  # remove1  $\psi$  ( $\varphi$  #  $\Phi$ )))) =
        (?m  $\circ$  image-mset ((#)  $\varphi$ )  $\circ$  image-mset ((#)  $\psi$ )) (mset ? $\Sigma$ ) +
        (?m  $\circ$  (image-mset (@) ( $\mathfrak{V}$   $\Delta'$ ))  $\circ$  image-mset ((#)  $\psi$ )) (mset
?Σ) +
        (?m  $\circ$  (image-mset ((#)  $\varphi$ )  $\circ$  image-mset (@) ( $\mathfrak{V}$   $\Delta$ ))) (mset
?Σ) +
        (?m  $\circ$  (image-mset (@) ( $\mathfrak{V}$   $\Delta'$ ))  $\circ$  image-mset (@) ( $\mathfrak{V}$   $\Delta$ )))
(mset ?Σ)
        by (simp add: image-mset-cons-homomorphism
          image-mset-append-homomorphism
          image-mset-add-collapse
          add-mset-commute
          add.commute)
      hence mset (map mset (?T $\Sigma$  ( $\psi$  # remove1  $\psi$  ( $\varphi$  #  $\Phi$ )))) =
        (?m  $\circ$  image-mset ((#)  $\varphi$ )) (mset ? $\Sigma'$ ) +
        (?m  $\circ$  (image-mset (@) ( $\mathfrak{V}$   $\Delta'$ ))) (mset ? $\Sigma'$ )
        using  $\langle \psi = (\Delta, \delta) \rangle$ 
        by (simp add: multiset.map-comp)
      hence mset (map mset (?T $\Sigma$  ( $\psi$  # remove1  $\psi$  ( $\varphi$  #  $\Phi$ )))) =
        image-mset ((+) {# $\varphi$ #}) (mset (map mset ? $\Sigma'$ )) +
        image-mset ((+) (mset ( $\mathfrak{V}$   $\Delta'$ ))) (mset (map mset ? $\Sigma'$ ))

```



by (simp add: image-mset-cons-homomorphism  
           image-mset-append-homomorphism)  
 hence mset (map mset (?T<sub>Σ</sub> (ψ # remove1 ψ (φ # Φ)))) =  
           image-mset ((+) {#φ#}) (mset (map mset (?T<sub>Σ</sub> Φ))) +  
           image-mset ((+) (mset (ℳ Δ'))) (mset (map mset (?T<sub>Σ</sub> Φ)))  
 using † by auto  
 hence mset (map mset (?T<sub>Σ</sub> (ψ # remove1 ψ (φ # Φ)))) =  
           (?m ∘ (image-mset ((#) φ))) (mset (?T<sub>Σ</sub> Φ)) +  
           (?m ∘ (image-mset ((@) (ℳ Δ')))) (mset (?T<sub>Σ</sub> Φ))  
 by (simp add: image-mset-cons-homomorphism  
           image-mset-append-homomorphism)  
 thus ?thesis using ⟨φ = (Δ', δ')⟩ by (simp add: multiset.map-comp)  
 qed  
 qed  
 hence image-mset mset (image-mset ((#) ψ) (mset (?T<sub>Σ</sub> (remove1 ψ  
 Φ)))) +  
           image-mset mset (image-mset ((@) (ℳ Δ)) (mset (?T<sub>Σ</sub> (remove1  
 ψ Φ))))  
           = image-mset mset (mset (?T<sub>Σ</sub> Φ))  
 by (simp add: ⟨ψ = (Δ, δ)⟩ multiset.map-comp)  
 hence  
           image-mset ((+) {# ψ #}) (image-mset mset (mset (?T<sub>Σ</sub> (remove1 ψ  
 Φ)))) +  
           image-mset ((+) (mset (ℳ Δ))) (image-mset mset (mset (?T<sub>Σ</sub> (remove1  
 ψ Φ))))  
           = image-mset mset (mset (?T<sub>Σ</sub> Φ))  
 by (simp add: image-mset-cons-homomorphism image-mset-append-homomorphism)  
 hence  
           image-mset ((+) {# ψ #}) (image-mset mset (mset (?T<sub>Σ</sub> Ψ<sub>0</sub>))) +  
           image-mset ((+) (mset (ℳ Δ))) (image-mset mset (mset (?T<sub>Σ</sub> Ψ<sub>0</sub>)))  
 = image-mset mset (mset (?T<sub>Σ</sub> Φ))  
 using Cons ⟨mset Ψ<sub>0</sub> = mset (remove1 ψ Φ)⟩  
 by fastforce  
 hence  
           image-mset mset (image-mset ((#) ψ) (mset (?T<sub>Σ</sub> Ψ<sub>0</sub>))) +  
           image-mset mset (image-mset ((@) (ℳ Δ)) (mset (?T<sub>Σ</sub> Ψ<sub>0</sub>)))  
 = image-mset mset (mset (?T<sub>Σ</sub> Φ))  
 by (simp add: image-mset-cons-homomorphism image-mset-append-homomorphism)  
 hence mset (map mset (?T<sub>Σ</sub> (ψ # Ψ<sub>0</sub>))) = mset (map mset (?T<sub>Σ</sub> Φ))  
 by (simp add: ⟨ψ = (Δ, δ)⟩ multiset.map-comp)  
 }  
 then show ?case by blast  
 qed  
 hence mset (map mset (?T<sub>Σ</sub> Ψ<sub>0</sub>)) = mset (map mset (?T<sub>Σ</sub> (?S Ψ<sub>0</sub>)))  
 using ⟨mset Ψ<sub>0</sub> = mset (?S Ψ<sub>0</sub>)⟩ by blast  
 hence mset (map (mset ∘ (map snd) ∘ remdups) (?T<sub>Σ</sub> Ψ<sub>0</sub>))  
           = mset (map (mset ∘ (map snd) ∘ remdups) (?T<sub>Σ</sub> (?S Ψ<sub>0</sub>)))  
 using mset-mset-map-snd-remdups by blast  
 hence mset (map mset ?Σ) = mset (map mset ?Σ')

```

    by (simp add: fun.map-comp)
  hence set (map mset ?Σ) = set (map mset ?Σ')
    using mset-eq-setD by blast
  hence  $\forall \sigma \in \text{set } ?\Sigma. \exists \sigma' \in \text{set } ?\Sigma'. \text{mset } \sigma = \text{mset } \sigma'$ 
    by fastforce
  hence  $\forall \sigma \in \text{set } ?\Sigma. \exists \sigma' \in \text{set } ?\Sigma'. \text{length } \sigma = \text{length } \sigma'$ 
    using mset-eq-length by blast
  have mset (?S Ψ0) ⊆# mset (ℳ Ξ)
    by (simp add: Ψ0(1))
  {
    fix n
    have  $\forall \Psi. \text{mset } \Psi \subseteq\# \text{mset } (\mathfrak{M} \Xi) \longrightarrow$ 
      sorted (map (- ?F) Ψ)  $\longrightarrow$ 
      length Ψ = n  $\longrightarrow$ 
      ( $\forall \sigma' \in \text{set } (\text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma \Psi)). \text{length } \sigma' + 1$ 
 $\geq n$ )
    proof (induct n)
      case 0
      then show ?case by simp
    next
      case (Suc n)
      {
        fix Ψ :: ('a list × 'a) list
        assume A: mset Ψ ⊆# mset (ℳ Ξ)
          and B: sorted (map (- ?F) Ψ)
          and C: length Ψ = n + 1
        obtain Δ δ where (Δ, δ) = hd Ψ
          using prod.collapse by blast
        let ?Ψ' = tl Ψ
        have mset ?Ψ' ⊆# mset (ℳ Ξ) using A
        by (induct Ψ, simp, simp, meson mset-subset-eq-insertD subset-mset-def)
        moreover
        have sorted (map (- ?F) (tl Ψ))
          using B
          by (simp add: map-tl sorted-tl)
        moreover have length ?Ψ' = n using C
          by simp
        ultimately have *:  $\forall \sigma' \in \text{set } (\text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_\Sigma \Psi')).$ 
 $\text{length } \sigma' + 1 \geq n$ 
          using Suc
          by blast
        from C have Ψ = (Δ, δ) # ?Ψ'
          by (metis ⟨(Δ, δ) = hd Ψ⟩
            One-nat-def
            add-is-0
            list.exhaust-sel
            list.size(3)
            nat.simps(3))
        have distinct ((Δ, δ) # ?Ψ')

```

```

using A ⟨Ψ = (Δ, δ) # ?Ψ'⟩
      MaxSAT-optimal-pre-witness-distinct
      mset-distinct-msub-down
by fastforce
hence set ((Δ, δ) # ?Ψ') ⊆ set (⋈ Ξ)
by (metis A ⟨Ψ = (Δ, δ) # ?Ψ'⟩
      Un-iff
      mset-le-perm-append
      perm-set-eq set-append
      subsetI)
hence ∀ (Δ', δ') ∈ set ?Ψ'. (Δ, δ) ≠ (Δ', δ')
      ∀ (Δ', δ') ∈ set (⋈ Ξ). ((Δ, δ) ≠ (Δ', δ')) ⟶ (length Δ ≠ length
Δ')
      set ?Ψ' ⊆ set (⋈ Ξ)
using MaxSAT-optimal-pre-witness-length-iff-eq [where Ψ=Ξ]
      ⟨distinct ((Δ, δ) # ?Ψ')⟩
by auto
hence ∀ (Δ', δ') ∈ set ?Ψ'. length Δ ≠ length Δ'
      sorted (map (- ?F) ((Δ, δ) # ?Ψ'))
using B ⟨Ψ = (Δ, δ) # ?Ψ'⟩
by (fastforce, auto)
hence ∀ (Δ', δ') ∈ set ?Ψ'. length Δ > length Δ'
by fastforce
{
  fix σ' :: 'a list
  assume σ' ∈ set (map (map snd ∘ remdups) (?TΣ Ψ))
  hence σ' ∈ set (map (map snd ∘ remdups) (?TΣ ((Δ, δ) # ?Ψ')))
    using ⟨Ψ = (Δ, δ) # ?Ψ'⟩
    by simp
  from this obtain ψ where ψ:
    ψ ∈ set (?TΣ ?Ψ')
    σ' = (map snd ∘ remdups ∘ (#) (Δ, δ)) ψ ∨
    σ' = (map snd ∘ remdups ∘ (@) (⋈ Δ)) ψ
    by fastforce
  hence length σ' ≥ n
  proof (cases σ' = (map snd ∘ remdups ∘ (#) (Δ, δ)) ψ)
    case True
    {
      fix Ψ :: ('a list × 'a) list
      fix n :: nat
      assume ∀ (Δ, δ) ∈ set Ψ. n > length Δ
      hence ∀ σ ∈ set (?TΣ Ψ). ∀ (Δ, δ) ∈ set σ. n > length Δ
      proof (induct Ψ)
        case Nil
        then show ?case by simp
      next
        case (Cons ψ Ψ)
        obtain Δ δ where ψ = (Δ, δ)
        by fastforce
    }
}

```

```

hence  $n > \text{length } \Delta$  using Cons.prems by fastforce
have  $0: \forall \sigma \in \text{set } (?T_{\Sigma} \Psi). \forall (\Delta', \delta') \in \text{set } \sigma. n > \text{length } \Delta'$ 
using Cons by simp
{
  fix  $\sigma :: ('a \text{ list} \times 'a) \text{ list}$ 
  fix  $\psi' :: 'a \text{ list} \times 'a$ 
  assume  $1: \sigma \in \text{set } (?T_{\Sigma} (\psi \# \Psi))$ 
    and  $2: \psi' \in \text{set } \sigma$ 
  obtain  $\Delta' \delta'$  where  $\psi' = (\Delta', \delta')$ 
    by fastforce
  have  $3: \sigma \in (\#) (\Delta, \delta) \text{ ' set } (?T_{\Sigma} \Psi) \vee \sigma \in (@) (\mathfrak{A} \Delta) \text{ ' set}$ 
    ( $?T_{\Sigma} \Psi$ )

    using  $1 \langle \psi = (\Delta, \delta) \rangle$  by simp
  have  $n > \text{length } \Delta'$ 
  proof (cases  $\sigma \in (\#) (\Delta, \delta) \text{ ' set } (?T_{\Sigma} \Psi)$ )
    case True
      from this obtain  $\sigma'$  where
         $\text{set } \sigma = \text{insert } (\Delta, \delta) (\text{set } \sigma')$ 
         $\sigma' \in \text{set } (?T_{\Sigma} \Psi)$ 
        by auto
      then show ?thesis
        using  $0 \langle \psi' \in \text{set } \sigma \rangle \langle \psi' = (\Delta', \delta') \rangle \langle n > \text{length } \Delta \rangle$ 
        by auto
    next
      case False
      from this and  $3$  obtain  $\sigma'$  where  $\sigma'$ :
         $\text{set } \sigma = \text{set } (\mathfrak{A} \Delta) \cup (\text{set } \sigma')$ 
         $\sigma' \in \text{set } (?T_{\Sigma} \Psi)$ 
        by auto
      have  $\forall (\Delta', \delta') \in \text{set } (\mathfrak{A} \Delta). \text{length } \Delta > \text{length } \Delta'$ 
        by (metis (mono-tags, lifting))
          case-prodI2
          MaxSAT-optimal-pre-witness-nonelement
          not-le
      hence  $\forall (\Delta', \delta') \in \text{set } (\mathfrak{A} \Delta). n > \text{length } \Delta'$ 
        using  $\langle n > \text{length } \Delta \rangle$  by auto
      then show ?thesis using  $0 \sigma' \langle \psi' \in \text{set } \sigma \rangle \langle \psi' = (\Delta', \delta') \rangle$  by
        (fastforce)
    qed
  hence  $n > \text{length } (\text{fst } \psi')$  using  $\langle \psi' = (\Delta', \delta') \rangle$  by fastforce
}
then show ?case by fastforce
qed
}
hence  $\forall \sigma \in \text{set } (?T_{\Sigma} ?\Psi'). \forall (\Delta', \delta') \in \text{set } \sigma. \text{length } \Delta > \text{length } \Delta'$ 
using  $\langle \forall (\Delta', \delta') \in \text{set } ?\Psi'. \text{length } \Delta > \text{length } \Delta' \rangle$ 
by blast
then show ?thesis using True  $\star \psi(1)$  by fastforce
next

```

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case False
have  $\forall (\Delta', \delta') \in \text{set } ?\Psi'. \text{length } \Delta \geq \text{length } \Delta'$ 
  using  $\langle \forall (\Delta', \delta') \in \text{set } ?\Psi'. \text{length } \Delta > \text{length } \Delta' \rangle$ 
  by auto
hence  $\forall (\Delta', \delta') \in \text{set } \Psi. \text{length } \Delta \geq \text{length } \Delta'$ 
  using  $\langle \Psi = (\Delta, \delta) \# ?\Psi' \rangle$ 
  by (metis case-prodI2 eq-iff prod.sel(1) set-ConsD)
hence  $\text{length } \Delta + 1 \geq \text{length } \Psi$ 
  using A MaxSAT-optimal-pre-witness-pigeon-hole
  by fastforce
hence  $\text{length } \Delta \geq n$ 
  using C
  by simp
have  $\text{length } \Delta = \text{length } (\mathfrak{V} \Delta)$ 
  by (induct  $\Delta$ , simp+)
hence  $\text{length } (\text{remdups } (\mathfrak{V} \Delta)) = \text{length } (\mathfrak{V} \Delta)$ 
  by (simp add: MaxSAT-optimal-pre-witness-distinct)
hence  $\text{length } (\text{remdups } (\mathfrak{V} \Delta)) \geq n$ 
  using  $\langle \text{length } \Delta = \text{length } (\mathfrak{V} \Delta) \rangle \langle n \leq \text{length } \Delta \rangle$ 
  by linarith
have  $\text{mset } (\text{remdups } (\mathfrak{V} \Delta @ \psi)) = \text{mset } (\text{remdups } (\psi @ \mathfrak{V} \Delta))$ 
  by (simp add: mset-remdups)
hence  $\text{length } (\text{remdups } (\mathfrak{V} \Delta @ \psi)) \geq \text{length } (\text{remdups } (\mathfrak{V} \Delta))$ 
  by (metis le-cases length-sub-mset mset-remdups-append-msub
size-mset)
hence  $\text{length } (\text{remdups } (\mathfrak{V} \Delta @ \psi)) \geq n$ 
  using  $\langle n \leq \text{length } (\text{remdups } (\mathfrak{V} \Delta)) \rangle$  dual-order.trans by blast
thus ?thesis using False  $\psi(2)$ 
  by simp
qed
}
hence  $\forall \sigma' \in \text{set } (\text{map } (\text{map } \text{snd} \circ \text{remdups}) (?T_{\Sigma} \Psi)). \text{length } \sigma' \geq n$ 
  by blast
}
then show ?case by fastforce
qed
}
hence  $\forall \sigma' \in \text{set } ?\Sigma'. \text{length } \sigma' + 1 \geq \text{length } (?S \Psi_0)$ 
  using  $\langle \text{mset } (?S \Psi_0) \subseteq \# \text{mset } (\mathfrak{V} \Xi) \rangle$ 
  by fastforce
hence  $\forall \sigma' \in \text{set } ?\Sigma'. \text{length } \sigma' + 1 \geq \text{length } \Psi_0$  by simp
hence  $\forall \sigma \in \text{set } ?\Sigma. \text{length } \sigma + 1 \geq \text{length } \Psi_0$ 
  using  $\langle \forall \sigma \in \text{set } ?\Sigma. \exists \sigma' \in \text{set } ?\Sigma'. \text{length } \sigma = \text{length } \sigma' \rangle$ 
  by fastforce
thus ?thesis using  $\Psi_0$  by fastforce
qed
have III:  $\forall \sigma \in \text{set } ?\Sigma. \text{mset } \sigma \subseteq \# \text{mset } \Xi$ 
proof -
  have  $\text{remdups } (\mathfrak{V} \Xi) = \mathfrak{V} \Xi$ 

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    by (simp add: MaxSAT-optimal-pre-witness-distinct distinct-remdups-id)
  from  $\Psi_0(1)$  have set  $\Psi_0 \subseteq \text{set } (\mathfrak{V} \Xi)$ 
    by (metis (no-types, lifting) ‹remdups  $(\mathfrak{V} \Xi) = \mathfrak{V} \Xi$ ›
        mset-remdups-set-sub-iff
        mset-remdups-subset-eq
        subset-mset.dual-order.trans)
  hence  $\forall \sigma \in \text{set } (?T_\Sigma \Psi_0). \text{set } \sigma \subseteq \text{set } (\mathfrak{V} \Xi)$ 
  proof (induct  $\Psi_0$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\psi \Psi_0$ )
    hence  $\forall \sigma \in \text{set } (?T_\Sigma \Psi_0). \text{set } \sigma \subseteq \text{set } (\mathfrak{V} \Xi)$  by auto
    obtain  $\Delta \delta$  where  $\psi = (\Delta, \delta)$  by fastforce
    hence  $(\Delta, \delta) \in \text{set } (\mathfrak{V} \Xi)$  using Cons by simp
    {
      fix  $\sigma :: ('a \text{ list} \times 'a \text{ list})$ 
      assume  $\star: \sigma \in (\#) (\Delta, \delta) \text{ ' set } (?T_\Sigma \Psi_0) \cup (@) (\mathfrak{V} \Delta) \text{ ' set } (?T_\Sigma \Psi_0)$ 
      have  $\text{set } \sigma \subseteq \text{set } (\mathfrak{V} \Xi)$ 
      proof (cases  $\sigma \in (\#) (\Delta, \delta) \text{ ' set } (?T_\Sigma \Psi_0)$ )
        case True
        then show ?thesis
          using ‹ $\forall \sigma \in \text{set } (?T_\Sigma \Psi_0). \text{set } \sigma \subseteq \text{set } (\mathfrak{V} \Xi)$ › ‹ $(\Delta, \delta) \in \text{set } (\mathfrak{V} \Xi)$ ›
          by fastforce
        case False
        hence  $\sigma \in (@) (\mathfrak{V} \Delta) \text{ ' set } (?T_\Sigma \Psi_0)$  using  $\star$  by simp
        moreover have  $\text{set } (\mathfrak{V} \Delta) \subseteq \text{set } (\mathfrak{V} \Xi)$ 
          using MaxSAT-optimal-pre-witness-element-inclusion ‹ $(\Delta, \delta) \in \text{set } (\mathfrak{V} \Xi)$ ›
          by fastforce
        ultimately show ?thesis
          using ‹ $\forall \sigma \in \text{set } (?T_\Sigma \Psi_0). \text{set } \sigma \subseteq \text{set } (\mathfrak{V} \Xi)$ ›
          by force
      }
    qed
  }
  hence  $\forall \sigma \in (\#) (\Delta, \delta) \text{ ' set } (?T_\Sigma \Psi_0) \cup (@) (\mathfrak{V} \Delta) \text{ ' set } (?T_\Sigma \Psi_0). \text{set } \sigma$ 
 $\subseteq \text{set } (\mathfrak{V} \Xi)$ 
    by auto
    thus ?case using ‹ $\psi = (\Delta, \delta)$ › by simp
  qed
  hence  $\forall \sigma \in \text{set } (?T_\Sigma \Psi_0). \text{mset } (\text{remdups } \sigma) \subseteq\# \text{mset } (\text{remdups } (\mathfrak{V} \Xi))$ 
    using mset-remdups-set-sub-iff by blast
  hence  $\forall \sigma \in \text{set } ?\Sigma. \text{mset } \sigma \subseteq\# \text{mset } (\text{map snd } (\mathfrak{V} \Xi))$ 
    using map-monotonic ‹remdups  $(\mathfrak{V} \Xi) = \mathfrak{V} \Xi$ ›
    by auto
  moreover have  $\text{map snd } (\mathfrak{V} \Xi) = \Xi$  by (induct  $\Xi$ , simp+)
  ultimately show ?thesis by simp
qed

```

**show** *?thesis* **using** *I II III* **by** *fastforce*  
**qed**  
**from** *this* **obtain**  $\Sigma_0$  **where**  $\Sigma_0$ :  
 $\vdash (\Psi \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap \Sigma_0) \rightarrow \varphi)$   
 $\forall \sigma \in \text{set } \Sigma_0. \text{mset } \sigma \subseteq\# \text{mset } \Xi \wedge \text{length } \sigma + 1 \geq \text{length } \Psi$   
**by** *blast*  
**moreover**  
**have**  $(\Phi @ \Psi) \rightarrow \varphi = \Phi \rightarrow (\Psi \rightarrow \varphi)$  **by** (*induct*  $\Phi$ , *simp+*)  
**hence**  $\vdash ((\Phi @ \Psi) \rightarrow \varphi) \leftrightarrow (\sqcap \Phi \rightarrow (\Psi \rightarrow \varphi))$   
**by** (*simp add: list-curry-uncurry*)  
**moreover** **have**  $\vdash (\Psi \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap \Sigma_0) \rightarrow \varphi)$   
 $\rightarrow (\Phi @ \Psi) \rightarrow \varphi \leftrightarrow (\sqcap \Phi \rightarrow \Psi \rightarrow \varphi)$   
 $\rightarrow (\Phi @ \Psi) \rightarrow \varphi \leftrightarrow ((\sqcap \Phi \sqcap \bigsqcup (\text{map } \sqcap \Sigma_0)) \rightarrow \varphi)$   
**proof** –  
**let**  $?\varphi = \langle \Psi \rightarrow \varphi \rangle \leftrightarrow (\langle \bigsqcup (\text{map } \sqcap \Sigma_0) \rangle \rightarrow \langle \varphi \rangle)$   
 $\rightarrow \langle (\Phi @ \Psi) \rightarrow \varphi \rangle \leftrightarrow (\langle \sqcap \Phi \rangle \rightarrow \langle \Psi \rightarrow \varphi \rangle)$   
 $\rightarrow \langle (\Phi @ \Psi) \rightarrow \varphi \rangle \leftrightarrow (\langle \sqcap \Phi \rangle \sqcap \langle \bigsqcup (\text{map } \sqcap \Sigma_0) \rangle) \rightarrow \langle \varphi \rangle$   
**have**  $\forall \mathfrak{M}. \mathfrak{M} \models_{prop} ?\varphi$  **by** *fastforce*  
**hence**  $\vdash (\langle ?\varphi \rangle)$  **using** *propositional-semantic* **by** *blast*  
**thus** *?thesis* **by** *simp*  
**qed**  
**moreover**  
**let**  $?\Sigma = \text{map } ((@) \Phi) \Sigma_0$   
**have**  $\forall \varphi \psi \chi. \vdash (\varphi \rightarrow \psi) \rightarrow \chi \rightarrow \psi \vee \neg \vdash \chi \rightarrow \varphi$   
**by** (*meson modus-ponens flip-hypothetical-syllogism*)  
**hence**  $\vdash ((\sqcap \Phi \sqcap \bigsqcup (\text{map } \sqcap \Sigma_0)) \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap ?\Sigma) \rightarrow \varphi)$   
**using** *append-dnf-distribute biconditional-def* **by** *fastforce*  
**ultimately** **have**  $\vdash (\Phi @ \Psi) \rightarrow \varphi \leftrightarrow (\bigsqcup (\text{map } \sqcap ?\Sigma) \rightarrow \varphi)$   
**using** *modus-ponens biconditional-transitivity-rule*  
**by** *blast*  
**moreover**  
{  
**fix**  $\sigma$   
**assume**  $\sigma \in \text{set } ?\Sigma$   
**from** *this* **obtain**  $\sigma_0$  **where**  $\sigma_0: \sigma = \Phi @ \sigma_0 \sigma_0 \in \text{set } \Sigma_0$  **by** (*simp*, *blast*)  
**hence**  $\text{mset } \sigma_0 \subseteq\# \text{mset } \Xi$  **using**  $\Sigma_0(2)$  **by** *blast*  
**hence**  $\text{mset } \sigma \subseteq\# \text{mset } (\Phi @ \Xi)$  **using**  $\sigma_0(1)$  **by** *simp*  
**hence**  $\text{mset } \sigma \subseteq\# \text{mset } \Gamma$  **using** *assms(1) assms(2)*  
**by** (*simp, meson subset-mset.dual-order.trans subset-mset.le-diff-conv2*)  
**moreover**  
**have**  $\text{length } \sigma + 1 \geq \text{length } (\Phi @ \Psi)$  **using**  $\Sigma_0(2) \sigma_0$  **by** *simp*  
**ultimately** **have**  $\text{mset } \sigma \subseteq\# \text{mset } \Gamma$   $\text{length } \sigma + 1 \geq \text{length } (\Phi @ \Psi)$  **by** *auto*  
}  
**ultimately**  
**show** *?thesis* **by** *blast*  
**qed**  
**lemma** (in *classical-logic*) *relative-maximals-optimal-witness*:  
**assumes**  $\neg \vdash \varphi$

**shows**  $0 < (\|\Gamma\|_\varphi)$   
 $= (\exists \Sigma. \text{mset} (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$   
 $\text{map} (\text{uncurry } (\sqcup)) \Sigma \vdash \varphi \wedge$   
 $1 + (\|\text{map} (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma\|_\varphi) = \|\Gamma\|_\varphi)$

**proof** (*rule iffI*)  
**assume**  $0 < \|\Gamma\|_\varphi$   
**from this obtain**  $\Xi$  **where**  $\Xi: \Xi \in \mathcal{M} \Gamma \varphi$  *length*  $\Xi < \text{length } \Gamma$   
**using**  $\langle \neg \vdash \varphi \rangle$   
*complement-relative-MaxSAT-def*  
*relative-MaxSAT-intro*  
*relative-maximals-existence*  
**by** *fastforce*  
**from this obtain**  $\psi$  **where**  $\psi: \psi \in \text{set} (\Gamma \ominus \Xi)$   
**by** (*metis*  $\langle 0 < \|\Gamma\|_\varphi \rangle$ )  
*less-not-refl*  
*list.exhaust*  
*list.set-intros(1)*  
*list.size(3)*  
*complement-relative-MaxSAT-intro*  
**let**  $?\Sigma = \mathfrak{W} \varphi (\psi \# \Xi)$   
**let**  $?\Sigma_A = \mathfrak{W}_\sqcup \varphi (\psi \# \Xi)$   
**let**  $?\Sigma_B = \mathfrak{W}_\rightarrow \varphi (\psi \# \Xi)$   
**have**  $\diamond: \text{mset} (\psi \# \Xi) \subseteq\# \text{mset } \Gamma$   
 $\psi \# \Xi \vdash \varphi$   
**using**  $\Xi(1) \psi$   
*relative-maximals-def*  
*list-deduction-theorem*  
*relative-maximals-complement-deduction*  
*mset-list-subtract-elem-cons-msub [where  $\Xi = \Xi$ ]*  
**by** *blast+*  
**moreover have**  $\text{map snd } ?\Sigma = \psi \# \Xi$  **by** (*induct*  $\Xi$ , *simp+*)  
**ultimately have**  $?\Sigma_A \vdash \varphi$   
 $\text{mset} (\text{map snd } ?\Sigma) \subseteq\# \text{mset } \Gamma$   
**using** *MaxSAT-optimal-witness-deduction*  
*list-deduction-def weak-biconditional-weaken*  
**by** (*metis+*)  
**moreover**  
{  
**let**  $?\Gamma' = ?\Sigma_B @ \Gamma \ominus \text{map snd } ?\Sigma$   
**have**  $A: \text{length } ?\Sigma_B = 1 + \text{length } \Xi$   
**by** (*induct*  $\Xi$ , *simp+*)  
**have**  $B: ?\Sigma_B \in \mathcal{M} ?\Gamma' \varphi$   
**proof** –  
**have**  $\neg ?\Sigma_B \vdash \varphi$   
**by** (*metis* (*no-types*, *lifting*))  
 $\Xi(1) \langle ?\Sigma_A \vdash \varphi \rangle$   
*modus-ponens list-deduction-def*  
*optimal-witness-split-identity*  
*relative-maximals-def*



*mem-Collect-eq*)

**moreover** have  $mset \ ?\Sigma_B \subseteq\# mset \ ?\Gamma'$   
 by *simp*

**hence**  $\forall \Psi. mset \ \Psi \subseteq\# mset \ ?\Gamma' \longrightarrow \neg \Psi \vdash \varphi \longrightarrow length \ \Psi \leq length \ ?\Sigma_B$

**proof** –

have  $\forall \Psi \in \mathcal{M} \ ?\Gamma' \varphi. length \ \Psi = length \ ?\Sigma_B$

**proof** (*rule ccontr*)

assume  $\neg (\forall \Psi \in \mathcal{M} \ ?\Gamma' \varphi. length \ \Psi = length \ ?\Sigma_B)$

**from this obtain**  $\Psi$  **where**

$\Psi: \Psi \in \mathcal{M} \ ?\Gamma' \varphi$   
 $length \ \Psi \neq length \ ?\Sigma_B$

by *blast*

have  $length \ \Psi \geq length \ ?\Sigma_B$

using  $\Psi(1)$   
 $\langle \neg \ ?\Sigma_B \vdash \varphi \rangle$   
 $\langle mset \ ?\Sigma_B \subseteq\# mset \ ?\Gamma' \rangle$

**unfolding** *relative-maximals-def*

by *blast*

**hence**  $length \ \Psi > length \ ?\Sigma_B$

using  $\Psi(2)$

by *linarith*

**have**  $length \ \Psi = length \ (\Psi \ominus \ ?\Sigma_B) + length \ (\Psi \cap \ ?\Sigma_B)$   
 (is  $length \ \Psi = length \ ?A + length \ ?B$ )

by (*metis (no-types, lifting)*)  
 $length-append$   
 $list-diff-intersect-comp$   
 $mset-append$   
 $mset-eq-length$ )

{

fix  $\sigma$

assume  $mset \ \sigma \subseteq\# mset \ \Gamma$   
 $length \ \sigma + 1 \geq length \ (?A @ \ ?B)$

**hence**  $length \ \sigma + 1 \geq length \ \Psi$   
 using  $\langle length \ \Psi = length \ ?A + length \ ?B \rangle$

by *simp*

**hence**  $length \ \sigma + 1 > length \ ?\Sigma_B$   
 using  $\langle length \ \Psi > length \ ?\Sigma_B \rangle$  by *linarith*

**hence**  $length \ \sigma + 1 > length \ \Xi + 1$   
 using  $A$  by *simp*

**hence**  $length \ \sigma > length \ \Xi$  by *linarith*

have  $\sigma \vdash \varphi$

**proof** (*rule ccontr*)

assume  $\neg \sigma \vdash \varphi$

**hence**  $length \ \sigma \leq length \ \Xi$   
 using  $\langle mset \ \sigma \subseteq\# mset \ \Gamma \rangle \ \Xi(1)$

**unfolding** *relative-maximals-def*

by *blast*

**thus** *False* using  $\langle length \ \sigma > length \ \Xi \rangle$  by *linarith*

qed

```

}
moreover
have mset  $\Psi \subseteq\# \text{mset } ?\Gamma'$ 
   $\neg \Psi \vdash \varphi$ 
   $\forall \Phi. \text{mset } \Phi \subseteq\# \text{mset } ?\Gamma' \wedge \neg \Phi \vdash \varphi \longrightarrow \text{length } \Phi \leq \text{length } \Psi$ 
  using  $\Psi(1)$  relative-maximals-def by blast+
hence mset  $?A \subseteq\# \text{mset } (\Gamma \ominus \text{map snd } ?\Sigma)$ 
  by (simp add: add.commute subset-eq-diff-conv)
hence mset  $?A \subseteq\# \text{mset } (\Gamma \ominus (\psi \# \Xi))$ 
  using  $\langle \text{map snd } ?\Sigma = \psi \# \Xi \rangle$  by metis
moreover
have mset  $?B \subseteq\# \text{mset } (\mathfrak{W} \rightarrow \varphi (\psi \# \Xi))$ 
  using list-intersect-right-project by blast
ultimately obtain  $\Sigma$  where  $\Sigma: \vdash ((?A @ ?B) \rightarrow \varphi) \leftrightarrow (\bigsqcup (\text{map } \sqcap \Sigma)$ 
 $\rightarrow \varphi)$ 
   $\forall \sigma \in \text{set } \Sigma. \sigma \vdash \varphi$ 
  using  $\diamond$  optimal-witness-list-intersect-biconditional
  by metis
hence  $\vdash \bigsqcup (\text{map } \sqcap \Sigma) \rightarrow \varphi$ 
  using weak-disj-of-conj-equiv by blast
hence  $?A @ ?B \vdash \varphi$ 
  using  $\Sigma(1)$  modus-ponens list-deduction-def weak-biconditional-weaken
  by blast
moreover have  $\text{set } (?A @ ?B) = \text{set } \Psi$ 
  using list-diff-intersect-comp union-code set-mset-mset by metis
hence  $?A @ ?B \vdash \varphi = \Psi \vdash \varphi$ 
  using list-deduction-monotonic by blast
ultimately have  $\Psi \vdash \varphi$  by metis
thus False using  $\Psi(1)$  unfolding relative-maximals-def by blast
qed
moreover have  $\exists \Psi. \Psi \in \mathcal{M} \text{ } ?\Gamma' \varphi$ 
  using assms relative-maximals-existence by blast
ultimately show ?thesis
  using relative-maximals-def
  by fastforce
qed
ultimately show ?thesis
  unfolding relative-maximals-def
  by fastforce
qed
have  $C: \forall \Xi \Gamma \varphi. \Xi \in \mathcal{M} \Gamma \varphi \longrightarrow \text{length } \Xi = |\Gamma|_{\varphi}$ 
  using relative-MaxSAT-intro by blast
then have  $D: \text{length } \Xi = |\Gamma|_{\varphi}$ 
  using  $\langle \Xi \in \mathcal{M} \Gamma \varphi \rangle$  by blast
have
   $\forall (\Sigma :: 'a \text{ list}) \Gamma n. (\neg \text{mset } \Sigma \subseteq\# \text{mset } \Gamma \vee \text{length } (\Gamma \ominus \Sigma) \neq n) \vee \text{length } \Gamma$ 
 $= n + \text{length } \Sigma$ 
  using list-subtract-msub-eq by blast
then have  $E: \text{length } \Gamma = \text{length } (\Gamma \ominus \text{map snd } (\mathfrak{W} \varphi (\psi \# \Xi))) + \text{length } (\psi$ 

```

```

#  $\Xi$ )
  using  $\langle \text{map snd } (\mathfrak{W} \varphi (\psi \# \Xi)) = \psi \# \Xi \rangle \langle \text{mset } (\psi \# \Xi) \subseteq\# \text{mset } \Gamma \rangle$  by
presburger
  have  $1 + \text{length } \Xi = | \mathfrak{W}_{\rightarrow} \varphi (\psi \# \Xi) @ \Gamma \ominus \text{map snd } (\mathfrak{W} \varphi (\psi \# \Xi)) |_{\varphi}$ 
  using  $C B A$  by presburger
  hence  $1 + (\| \text{map } (\text{uncurry } (\rightarrow)) \text{ ?}\Sigma @ \Gamma \ominus \text{map snd } \text{ ?}\Sigma \|_{\varphi}) = \| \Gamma \|_{\varphi}$ 
  using  $D E \langle \text{map snd } (\mathfrak{W} \varphi (\psi \# \Xi)) = \psi \# \Xi \rangle$  complement-relative-MaxSAT-def
by force
}
ultimately
show  $\exists \Sigma. \text{mset } (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$ 
   $\text{map } (\text{uncurry } (\sqcup)) \Sigma \vdash \varphi \wedge$ 
   $1 + (\| \text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \|_{\varphi}) = \| \Gamma \|_{\varphi}$ 
by metis
next
assume  $\exists \Sigma. \text{mset } (\text{map snd } \Sigma) \subseteq\# \text{mset } \Gamma \wedge$ 
   $\text{map } (\text{uncurry } (\sqcup)) \Sigma \vdash \varphi \wedge$ 
   $1 + (\| \text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma \|_{\varphi}) = \| \Gamma \|_{\varphi}$ 
thus  $0 < \| \Gamma \|_{\varphi}$ 
by auto
qed

```

```

primrec (in implication-logic)
MaxSAT-witness :: ('a × 'a) list ⇒ 'a list ⇒ ('a × 'a) list (⟨ $\mathfrak{U}$ ⟩)
where
   $\mathfrak{U} - [] = []$ 
  |  $\mathfrak{U} \Sigma (\xi \# \Xi) = (\text{case find } (\lambda \sigma. \xi = \text{snd } \sigma) \Sigma \text{ of}$ 
     $\text{None} \Rightarrow \mathfrak{U} \Sigma \Xi$ 
    |  $\text{Some } \sigma \Rightarrow \sigma \# (\mathfrak{U} (\text{remove1 } \sigma \Sigma) \Xi))$ 

```

```

lemma (in implication-logic) MaxSAT-witness-right-msub:
   $\text{mset } (\text{map snd } (\mathfrak{U} \Sigma \Xi)) \subseteq\# \text{mset } \Xi$ 

```

```

proof -
have  $\forall \Sigma. \text{mset } (\text{map snd } (\mathfrak{U} \Sigma \Xi)) \subseteq\# \text{mset } \Xi$ 
proof (induct  $\Xi$ )
  case Nil
  then show ?case by simp
next
  case (Cons  $\xi \Xi$ )
  {
  fix  $\Sigma$ 
  have  $\text{mset } (\text{map snd } (\mathfrak{U} \Sigma (\xi \# \Xi))) \subseteq\# \text{mset } (\xi \# \Xi)$ 
  proof (cases find  $(\lambda \sigma. \xi = \text{snd } \sigma) \Sigma$ )
    case None
    then show ?thesis
    by (simp, metis Cons.hyps
      add-mset-add-single
      mset-map mset-subset-eq-add-left subset-mset.order-trans)

```

```

next
  case (Some  $\sigma$ )
  note  $\sigma = this$ 
  hence  $\xi = snd \sigma$ 
    by (meson find-Some-predicate)
  moreover
  have  $\sigma \in set \Sigma$ 
  using  $\sigma$ 
  proof (induct  $\Sigma$ )
    case Nil
    then show ?case by simp
  next
    case (Cons  $\sigma' \Sigma$ )
    then show ?case
      by (cases  $\xi = snd \sigma', simp+$ )
  qed
  ultimately show ?thesis using  $\sigma Cons.hyps$  by simp
qed
}
then show ?case by simp
qed
thus ?thesis by simp
qed

```

**lemma** (in *implication-logic*) *MaxSAT-witness-left-msub*:

$mset (\mathcal{U} \Sigma \Xi) \subseteq\# mset \Sigma$

**proof** –

have  $\forall \Sigma. mset (\mathcal{U} \Sigma \Xi) \subseteq\# mset \Sigma$

**proof** (induct  $\Xi$ )

case Nil

then show ?case by simp

next

case (Cons  $\xi \Xi$ )

{

fix  $\Sigma$

have  $mset (\mathcal{U} \Sigma (\xi \# \Xi)) \subseteq\# mset \Sigma$

**proof** (cases find  $(\lambda \sigma. \xi = snd \sigma) \Sigma$ )

case None

then show ?thesis using *Cons.hyps* by simp

next

case (Some  $\sigma$ )

note  $\sigma = this$

hence  $\sigma \in set \Sigma$

**proof** (induct  $\Sigma$ )

case Nil

then show ?case by simp

next

case (Cons  $\sigma' \Sigma$ )

then show ?case

```

      by (cases  $\xi = \text{snd } \sigma'$ , simp+)
    qed
    moreover from Cons.hyps have  $\text{mset } (\mathcal{U} (\text{remove1 } \sigma \Sigma) \Xi) \subseteq\# \text{mset}$ 
(remove1  $\sigma \Sigma$ )
      by blast
    hence  $\text{mset } (\mathcal{U} \Sigma (\xi \# \Xi)) \subseteq\# \text{mset } (\sigma \# \text{remove1 } \sigma \Sigma)$  using  $\sigma$  by simp
    ultimately show ?thesis by simp
  qed
}
then show ?case by simp
qed
thus ?thesis by simp
qed

```

**lemma** (in *implication-logic*) *MaxSAT-witness-right-projection:*

$\text{mset } (\text{map } \text{snd } (\mathcal{U} \Sigma \Xi)) = \text{mset } ((\text{map } \text{snd } \Sigma) \cap \Xi)$

**proof** –

**have**  $\forall \Sigma. \text{mset } (\text{map } \text{snd } (\mathcal{U} \Sigma \Xi)) = \text{mset } ((\text{map } \text{snd } \Sigma) \cap \Xi)$

**proof** (induct  $\Xi$ )

case *Nil*

then show ?case by simp

**next**

case (Cons  $\xi \Xi$ )

{

fix  $\Sigma$

**have**  $\text{mset } (\text{map } \text{snd } (\mathcal{U} \Sigma (\xi \# \Xi))) = \text{mset } (\text{map } \text{snd } \Sigma \cap \xi \# \Xi)$

**proof** (cases find  $(\lambda \sigma. \xi = \text{snd } \sigma) \Sigma$ )

case *None*

hence  $\xi \notin \text{set } (\text{map } \text{snd } \Sigma)$

**proof** (induct  $\Sigma$ )

case *Nil*

then show ?case by simp

**next**

case (Cons  $\sigma \Sigma$ )

**have** find  $(\lambda \sigma. \xi = \text{snd } \sigma) \Sigma = \text{None}$

$\xi \neq \text{snd } \sigma$

using Cons.prem

by (auto, metis Cons.prem find.simps(2) find-None-iff list.set-intros(1))

then show ?case using Cons.hyps by simp

qed

then show ?thesis using None Cons.hyps by simp

**next**

case (Some  $\sigma$ )

hence  $\sigma \in \text{set } \Sigma \wedge \xi = \text{snd } \sigma$

by (meson find-Some-predicate find-Some-set-membership)+

**moreover**

**from**  $\langle \sigma \in \text{set } \Sigma \rangle$  **have**  $\text{mset } \Sigma = \text{mset } (\sigma \# (\text{remove1 } \sigma \Sigma))$

by simp

**hence**  $\text{mset } (\text{map } \text{snd } \Sigma) = \text{mset } ((\text{snd } \sigma) \# (\text{remove1 } (\text{snd } \sigma) (\text{map } \text{snd}$

```

Σ)))
      mset (map snd Σ) = mset (map snd (σ # (remove1 σ Σ)))
      by (simp add: ⟨σ ∈ set Σ⟩, metis map-monotonic subset-mset.eq-iff)
      hence mset (map snd (remove1 σ Σ)) = mset (remove1 (snd σ) (map snd
Σ))
      by simp
      ultimately show ?thesis using Some Cons.hyps by simp
    qed
  }
  then show ?case by simp
  qed
  thus ?thesis by simp
qed

lemma (in classical-logic) witness-list-implication-rule:
  ⊢ (map (uncurry (⊔)) Σ :→ φ) → ∏ (map (λ (χ, ξ). (χ → ξ) → φ) Σ) → φ
proof (induct Σ)
  case Nil
  then show ?case using axiom-k by simp
next
  case (Cons σ Σ)
  let ?χ = fst σ
  let ?ξ = snd σ
  let ?ΣA = map (uncurry (⊔)) Σ
  let ?ΣB = map (λ (χ, ξ). (χ → ξ) → φ) Σ
  assume ⊢ ?ΣA :→ φ → ∏ ?ΣB → φ
  moreover have
    ⊢ (?ΣA :→ φ → ∏ ?ΣB → φ)
    → ((?χ ⊔ ?ξ) → ?ΣA :→ φ) → (((?χ → ?ξ) → φ) ∏ ∏ ?ΣB) → φ
  proof -
    let ?φ = (⟨?ΣA :→ φ⟩ → ⟨∏ ?ΣB⟩ → ⟨φ⟩)
    → (((⟨?χ⟩ ⊔ ⟨?ξ⟩) → ⟨?ΣA :→ φ⟩) → (((⟨?χ⟩ → ⟨?ξ⟩) → ⟨φ⟩) ∏
⟨∏ ?ΣB⟩) → ⟨φ⟩)
    have ∀ M. M ⊢prop ?φ by fastforce
    hence ⊢ (⟦ ?φ ⤵) using propositional-semantics by blast
    thus ?thesis by simp
  qed
  ultimately have ⊢ ((?χ ⊔ ?ξ) → ?ΣA :→ φ) → (((?χ → ?ξ) → φ) ∏ ∏ ?ΣB)
→ φ
  using modus-ponens by blast
  moreover
  have (λ σ. (fst σ → snd σ) → φ) = (λ (χ, ξ). (χ → ξ) → φ)
    uncurry (⊔) = (λ σ. fst σ ⊔ snd σ)
  by fastforce+
  hence (λ (χ, ξ). (χ → ξ) → φ) σ = (?χ → ?ξ) → φ
    uncurry (⊔) σ = ?χ ⊔ ?ξ
  by metis+
  ultimately show ?case by simp
qed

```

**lemma** (in *classical-logic*) *witness-relative-MaxSAT-increase*:

**assumes**  $\neg \vdash \varphi$   
**and**  $mset (map\ snd\ \Sigma) \subseteq\# mset\ \Gamma$   
**and**  $map (uncurry (\sqcup))\ \Sigma \vdash \varphi$   
**shows**  $(|\ \Gamma \ |_{\varphi}) < (|\ map (uncurry (\rightarrow))\ \Sigma @ \Gamma \ominus map\ snd\ \Sigma \ |_{\varphi})$

**proof** –

**from**  $\langle \neg \vdash \varphi \rangle$  **obtain**  $\Xi$  **where**  $\Xi: \Xi \in \mathcal{M}\ \Gamma\ \varphi$

**using** *relative-maximals-existence* **by** *blast*

**let**  $? \Sigma' = \Sigma \ominus \mathcal{U}\ \Sigma\ \Xi$

**let**  $? \Sigma \Xi' = map (uncurry (\sqcup)) (\mathcal{U}\ \Sigma\ \Xi) @ map (uncurry (\rightarrow)) (\mathcal{U}\ \Sigma\ \Xi)$

**have**  $mset\ \Sigma = mset (\mathcal{U}\ \Sigma\ \Xi @ ? \Sigma')$  **by** (*simp add: MaxSAT-witness-left-msub*)

**hence**  $set (map (uncurry (\sqcup))\ \Sigma) = set (map (uncurry (\sqcup)) (\mathcal{U}\ \Sigma\ \Xi @ ? \Sigma'))$

**by** (*metis mset-map mset-eq-setD*)

**hence**  $map (uncurry (\sqcup)) (\mathcal{U}\ \Sigma\ \Xi) @ ? \Sigma' \vdash \varphi$

**using** *list-deduction-monotonic assms(3)*

**by** *blast*

**hence**  $map (uncurry (\sqcup)) (\mathcal{U}\ \Sigma\ \Xi) @ map (uncurry (\sqcup))\ ? \Sigma' \vdash \varphi$  **by** *simp*

**moreover**

{

**fix**  $\Phi\ \Psi$

**have**  $((\Phi @ \Psi) \rightarrow \varphi) = (\Phi \rightarrow (\Psi \rightarrow \varphi))$

**by** (*induct*  $\Phi$ , *simp+*)

**hence**  $(\Phi @ \Psi) \vdash \varphi = \Phi \vdash (\Psi \rightarrow \varphi)$

**unfolding** *list-deduction-def*

**by** (*induct*  $\Phi$ , *simp+*)

}

**ultimately** **have**  $map (uncurry (\sqcup)) (\mathcal{U}\ \Sigma\ \Xi) \vdash map (uncurry (\sqcup))\ ? \Sigma' \rightarrow \varphi$

**by** *simp*

**moreover** **have**  $set (map (uncurry (\sqcup)) (\mathcal{U}\ \Sigma\ \Xi)) \subseteq set\ ? \Sigma \Xi'$

**by** *simp*

**ultimately** **have**  $? \Sigma \Xi' \vdash map (uncurry (\sqcup))\ ? \Sigma' \rightarrow \varphi$

**using** *list-deduction-monotonic* **by** *blast*

**hence**  $? \Sigma \Xi' \vdash \prod (map (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi)\ ? \Sigma') \rightarrow \varphi$

**using** *list-deduction-modus-ponens*

*list-deduction-weaken*

*witness-list-implication-rule*

**by** *blast*

**hence**  $? \Sigma \Xi' \$\vdash [\prod (map (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi)\ ? \Sigma') \rightarrow \varphi]$

**using** *measure-deduction-one-collapse* **by** *metis*

**hence**

$? \Sigma \Xi' @ (map\ snd (\mathcal{U}\ \Sigma\ \Xi)) \ominus (map\ snd (\mathcal{U}\ \Sigma\ \Xi))$

$\$ \vdash [\prod (map (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi)\ ? \Sigma') \rightarrow \varphi]$

**by** *simp*

**hence**  $map\ snd (\mathcal{U}\ \Sigma\ \Xi) \$\vdash [\prod (map (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi)\ ? \Sigma') \rightarrow \varphi]$

**using** *measure-witness-left-split* [**where**  $\Gamma = map\ snd (\mathcal{U}\ \Sigma\ \Xi)$

**and**  $\Sigma = \mathcal{U}\ \Sigma\ \Xi]$

**by** *fastforce*

**hence**  $map\ snd (\mathcal{U}\ \Sigma\ \Xi) \$\vdash [\prod (map (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi)\ ? \Sigma') \rightarrow \varphi]$

```

using MaxSAT-witness-right-projection by auto
hence  $\text{map snd } (\mathfrak{U} \Sigma \Xi) \vdash \prod (\text{map } (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi) \text{ ?}\Sigma') \rightarrow \varphi$ 
using measure-deduction-one-collapse by blast
hence  $\star$ :
   $\text{map snd } (\mathfrak{U} \Sigma \Xi) @ \Xi \ominus (\text{map snd } \Sigma) \vdash \prod (\text{map } (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi)$ 
 $\text{ ?}\Sigma') \rightarrow \varphi$ 
  (is  $\text{?}\Xi_0 \vdash -$ )
using list-deduction-monotonic
by (metis (no-types, lifting) append-Nil2
      measure-cancel
      measure-deduction.simps(1)
      measure-list-deduction-antitonic)
have  $\text{mset } \Xi = \text{mset } (\Xi \ominus (\text{map snd } \Sigma)) + \text{mset } (\Xi \cap (\text{map snd } \Sigma))$ 
using list-diff-intersect-comp by blast
hence  $\text{mset } \Xi = \text{mset } ((\text{map snd } \Sigma) \cap \Xi) + \text{mset } (\Xi \ominus (\text{map snd } \Sigma))$ 
by (metis subset-mset.inf-commute list-intersect-mset-homomorphism union-commute)
hence  $\text{mset } \Xi = \text{mset } (\text{map snd } (\mathfrak{U} \Sigma \Xi)) + \text{mset } (\Xi \ominus (\text{map snd } \Sigma))$ 
using MaxSAT-witness-right-projection by simp
hence  $\text{mset } \Xi = \text{mset } \text{?}\Xi_0$ 
by simp
hence  $\text{set } \Xi = \text{set } \text{?}\Xi_0$ 
by (metis mset-eq-setD)
have  $\neg \text{?}\Xi_0 \vdash \prod (\text{map } (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi) \text{ ?}\Sigma')$ 
proof (rule notI)
  assume  $\text{?}\Xi_0 \vdash \prod (\text{map } (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi) \text{ ?}\Sigma')$ 
  hence  $\text{?}\Xi_0 \vdash \varphi$ 
  using  $\star$  list-deduction-modus-ponens by blast
  hence  $\Xi \vdash \varphi$ 
  using list-deduction-monotonic  $\langle \text{set } \Xi = \text{set } \text{?}\Xi_0 \rangle$  by blast
thus False
using  $\Xi$  relative-maximals-def by blast
qed
moreover
have  $\text{mset } (\text{map snd } (\mathfrak{U} \Sigma \Xi)) \subseteq\# \text{mset } \text{?}\Xi_0$ 
   $\text{mset } (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{U} \Sigma \Xi) @ \text{?}\Xi_0 \ominus \text{map snd } (\mathfrak{U} \Sigma \Xi))$ 
   $= \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{U} \Sigma \Xi) @ \Xi \ominus (\text{map snd } \Sigma))$ 
  (is  $- = \text{mset } \text{?}\Xi_1$ )
by auto
hence  $\text{?}\Xi_1 \preceq \text{?}\Xi_0$ 
by (metis add.commute
      witness-stronger-theory
      add-diff-cancel-right'
      list-subtract.simps(1)
      list-subtract-mset-homomorphism
      list-diff-intersect-comp
      list-intersect-right-project
      msub-stronger-theory-intro
      stronger-theory-combine
      stronger-theory-empty-list-intro)

```



*self-append-conv*)

**ultimately have**  
 $\neg \text{?}\Xi_1 \vdash \sqcap (\text{map } (\lambda (\chi, \gamma). (\chi \rightarrow \gamma) \rightarrow \varphi) \text{?}\Sigma')$   
**using** *stronger-theory-deduction-monotonic* **by** *blast*

**from this obtain**  $\chi \ \gamma$  **where**  
 $(\chi, \gamma) \in \text{set } \text{?}\Sigma'$   
 $\neg (\chi \rightarrow \gamma) \# \text{?}\Xi_1 \vdash \varphi$   
**using** *list-deduction-theorem*  
**by** *fastforce*

**have**  $\text{mset } (\chi \rightarrow \gamma \# \text{?}\Xi_1) \subseteq \# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map } \text{snd } \Sigma)$

**proof** –  
**let**  $\text{?}A = \text{map } (\text{uncurry } (\rightarrow)) \Sigma$   
**let**  $\text{?}B = \text{map } (\text{uncurry } (\rightarrow)) (\mathfrak{U} \Sigma \Xi)$   
**have**  $(\chi, \gamma) \in (\text{set } \Sigma - \text{set } (\mathfrak{U} \Sigma \Xi))$   
**proof** –  
**from**  $(\chi, \gamma) \in \text{set } \text{?}\Sigma'$  **have**  $\gamma \in \# \text{mset } (\text{map } \text{snd } (\Sigma \ominus \mathfrak{U} \Sigma \Xi))$   
**by** *(metis set-mset-mset image-eqI set-map snd-conv)*  
**hence**  $\gamma \in \# \text{mset } (\text{map } \text{snd } \Sigma \ominus \text{map } \text{snd } (\mathfrak{U} \Sigma \Xi))$   
**by** *(metis MaxSAT-witness-left-msub map-list-subtract-mset-equivalence)*  
**hence**  $\gamma \in \# \text{mset } (\text{map } \text{snd } \Sigma \ominus (\text{map } \text{snd } \Sigma \cap \Xi))$   
**by** *(metis MaxSAT-witness-right-projection list-subtract-mset-homomorphism)*  
**hence**  $\gamma \in \# \text{mset } (\text{map } \text{snd } \Sigma \ominus \Xi)$   
**by** *(metis add-diff-cancel-right'*  
*list-subtract-mset-homomorphism*  
*list-diff-intersect-comp)*

**moreover from** *assms(2)* **have**  $\text{mset } (\text{map } \text{snd } \Sigma \ominus \Xi) \subseteq \# \text{mset } (\Gamma \ominus \Xi)$   
**by** *(simp, metis list-subtract-monotonic list-subtract-mset-homomorphism*  
*mset-map)*

**ultimately have**  $\gamma \in \# \text{mset } (\Gamma \ominus \Xi)$   
**by** *(simp add: mset-subset-eqD)*  
**hence**  $\gamma \in \text{set } (\Gamma \ominus \Xi)$   
**using** *set-mset-mset* **by** *fastforce*  
**hence**  $\gamma \in \text{set } \Gamma - \text{set } \Xi$   
**using**  $\Xi$  **by** *simp*  
**hence**  $\gamma \notin \text{set } \Xi$   
**by** *blast*  
**hence**  $\forall \Sigma. (\chi, \gamma) \notin \text{set } (\mathfrak{U} \Sigma \Xi)$   
**proof** *(induct*  $\Xi)$   
**case** *Nil*  
**then show** *?case* **by** *simp*

**next**  
**case** *(Cons*  $\xi \ \Xi)$   
**{**  
**fix**  $\Sigma$   
**have**  $(\chi, \gamma) \notin \text{set } (\mathfrak{U} \Sigma (\xi \# \Xi))$   
**proof** *(cases find*  $(\lambda \sigma. \xi = \text{snd } \sigma) \Sigma)$   
**case** *None*  
**then show** *?thesis* **using** *Cons* **by** *simp*  
**next**

```

      case (Some  $\sigma$ )
      moreover from this have  $\text{snd } \sigma = \xi$ 
        using find-Some-predicate by fastforce
      with Cons.premis have  $\sigma \neq (\chi, \gamma)$  by fastforce
      ultimately show ?thesis using Cons by simp
    qed
  }
  then show ?case by blast
qed
moreover from  $\langle (\chi, \gamma) \in \text{set } ?\Sigma' \rangle$  have  $(\chi, \gamma) \in \text{set } \Sigma$ 
  by (meson list-subtract-set-trivial-upper-bound subsetCE)
ultimately show ?thesis by fastforce
qed
with  $\langle (\chi, \gamma) \in \text{set } ?\Sigma' \rangle$  have  $\text{mset } ((\chi, \gamma) \# \mathfrak{A} \Sigma \Xi) \subseteq \# \text{mset } \Sigma$ 
  by (meson MaxSAT-witness-left-msub msub-list-subtract-elem-cons-msub)
hence  $\text{mset } (\chi \rightarrow \gamma \# ?B) \subseteq \# \text{mset } (\text{map } (\text{uncurry } (\rightarrow)) \Sigma)$ 
  by (metis (no-types, lifting)
     $\langle (\chi, \gamma) \in \text{set } ?\Sigma' \rangle$ 
    MaxSAT-witness-left-msub
    map-list-subtract-mset-equivalence
    map-monotonic
    mset-eq-setD msub-list-subtract-elem-cons-msub
    pair-imageI
    set-map
    uncurry-def)
moreover
have  $\text{mset } \Xi \subseteq \# \text{mset } \Gamma$ 
  using  $\Xi$  relative-maximals-def
  by blast
hence  $\text{mset } (\Xi \ominus (\text{map } \text{snd } \Sigma)) \subseteq \# \text{mset } (\Gamma \ominus (\text{map } \text{snd } \Sigma))$ 
  using list-subtract-monotonic by blast
ultimately show ?thesis
  using subset-mset.add-mono by fastforce
qed
moreover have  $\text{length } ?\Xi_1 = \text{length } ?\Xi_0$ 
  by simp
hence  $\text{length } ?\Xi_1 = \text{length } \Xi$ 
  using  $\langle \text{mset } \Xi = \text{mset } ?\Xi_0 \rangle$  mset-eq-length
  by metis
hence  $\text{length } ((\chi \rightarrow \gamma) \# ?\Xi_1) = \text{length } \Xi + 1$ 
  by simp
hence  $\text{length } ((\chi \rightarrow \gamma) \# ?\Xi_1) = (|\Gamma|_{|\varphi|} + 1)$ 
  using  $\Xi$ 
  by (simp add: relative-MaxSAT-intro)
moreover from  $\langle \neg \vdash \varphi \rangle$  obtain  $\Omega$  where  $\Omega: \Omega \in \mathcal{M} (\text{map } (\text{uncurry } (\rightarrow)) \Sigma @$ 
 $\Gamma \ominus \text{map } \text{snd } \Sigma) \varphi$ 
  using relative-maximals-existence by blast
ultimately have  $\text{length } \Omega \geq (|\Gamma|_{|\varphi|} + 1)$ 
  using relative-maximals-def

```

```

    by (metis (no-types, lifting)  $\langle \neg \chi \rightarrow \gamma \# \text{?}\Xi_1 \vdash \varphi \rangle$  mem-Collect-eq)
  thus ?thesis
    using  $\Omega$  relative-MaxSAT-intro by auto
qed

lemma (in classical-logic) relative-maximals-counting-deduction-lower-bound:
  assumes  $\neg \vdash \varphi$ 
  shows  $(\Gamma \# \vdash n \varphi) = (n \leq \|\Gamma\|_\varphi)$ 
proof -
  have  $\forall \Gamma. (\Gamma \# \vdash n \varphi) = (n \leq \|\Gamma\|_\varphi)$ 
  proof (induct n)
    case 0
    then show ?case by simp
  next
    case (Suc n)
    {
      fix  $\Gamma$ 
      assume  $\Gamma \# \vdash (\text{Suc } n) \varphi$ 
      from this obtain  $\Sigma$  where  $\Sigma$ :
        mset (map snd  $\Sigma$ )  $\subseteq \#$  mset  $\Gamma$ 
        map (uncurry ( $\sqcup$ ))  $\Sigma \vdash \varphi$ 
        map (uncurry ( $\rightarrow$ ))  $\Sigma @ \Gamma \ominus (\text{map snd } \Sigma) \# \vdash n \varphi$ 
      by fastforce
      let  $\text{?}\Gamma' = \text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus (\text{map snd } \Sigma)$ 
      have length  $\Gamma = \text{length } \text{?}\Gamma'$ 
      using  $\Sigma(1)$  list-subtract-msub-eq by fastforce
      hence  $(\|\Gamma\|_\varphi) > (\|\text{?}\Gamma'\|_\varphi)$ 
      by (metis  $\Sigma(1) \Sigma(2) \langle \neg \vdash \varphi \rangle$ 
          witness-relative-MaxSAT-increase
          length-MaxSAT-decomposition
          add-less-cancel-right
          nat-add-left-cancel-less)
      with  $\Sigma(3)$  Suc.hyps have  $\text{Suc } n \leq \|\Gamma\|_\varphi$ 
      by auto
    }
  moreover
  {
    fix  $\Gamma$ 
    assume  $\text{Suc } n \leq \|\Gamma\|_\varphi$ 
    from this obtain  $\Sigma$  where  $\Sigma$ :
      mset (map snd  $\Sigma$ )  $\subseteq \#$  mset  $\Gamma$ 
      map (uncurry ( $\sqcup$ ))  $\Sigma \vdash \varphi$ 
       $1 + (\|\text{map } (\text{uncurry } (\rightarrow)) \Sigma @ \Gamma \ominus \text{map snd } \Sigma\|_\varphi) = \|\Gamma\|_\varphi$ 
      (is  $1 + (\|\text{?}\Gamma'\|_\varphi) = \|\Gamma\|_\varphi$ )
    by (metis Suc-le-D assms relative-maximals-optimal-witness zero-less-Suc)
    have  $n \leq \|\text{?}\Gamma'\|_\varphi$ 
    using  $\Sigma(3) \langle \text{Suc } n \leq \|\Gamma\|_\varphi \rangle$  by linarith
    hence  $\text{?}\Gamma' \# \vdash n \varphi$  using Suc by blast
    hence  $\Gamma \# \vdash (\text{Suc } n) \varphi$  using  $\Sigma(1) \Sigma(2)$  by fastforce
  }

```

```

    }
    ultimately show ?case by metis
  qed
  thus ?thesis by auto
qed

```

As a brief aside, we may observe that  $\varphi$  is a tautology if and only if counting deduction can prove it for any given number of times. This follows immediately from  $\neg \vdash \varphi \implies \Gamma \# \vdash n \varphi = (n \leq \|\Gamma\|_\varphi)$ .

**lemma** (in *classical-logic*) *counting-deduction-tautology-equiv*:

$(\forall n. \Gamma \# \vdash n \varphi) = \vdash \varphi$

**proof** (*cases*  $\vdash \varphi$ )

**case** *True*

**then show** ?thesis

by (*simp add: counting-deduction-tautology-weaken*)

**next**

**case** *False*

**have**  $\neg \Gamma \# \vdash (1 + \text{length } \Gamma) \varphi$

**proof** (*rule notI*)

**assume**  $\Gamma \# \vdash (1 + \text{length } \Gamma) \varphi$

**hence**  $1 + \text{length } \Gamma \leq \|\Gamma\|_\varphi$

**using**  $\langle \neg \vdash \varphi \rangle$  *relative-maximals-counting-deduction-lower-bound* **by** *blast*

**hence**  $1 + \text{length } \Gamma \leq \text{length } \Gamma$

**using** *complement-relative-MaxSAT-def* **by** *fastforce*

**thus** *False* **by** *linarith*

**qed**

**then show** ?thesis

**using**  $\langle \neg \vdash \varphi \rangle$  **by** *blast*

**qed**

**theorem** (in *classical-logic*) *relative-maximals-max-counting-deduction*:

$\Gamma \# \vdash n \varphi = (\forall \Phi \in \mathcal{M} \Gamma \varphi. n \leq \text{length } (\Gamma \ominus \Phi))$

**proof** (*cases*  $\vdash \varphi$ )

**case** *True*

**from**  $\langle \vdash \varphi \rangle$  **have**  $\Gamma \# \vdash n \varphi$

**using** *counting-deduction-tautology-weaken*

**by** *blast*

**moreover from**  $\langle \vdash \varphi \rangle$  **have**  $\mathcal{M} \Gamma \varphi = \{\}$

**using** *relative-maximals-existence* **by** *auto*

**hence**  $\forall \Phi \in \mathcal{M} \Gamma \varphi. n \leq \text{length } (\Gamma \ominus \Phi)$  **by** *blast*

**ultimately show** ?thesis **by** *meson*

**next**

**case** *False*

**from**  $\langle \neg \vdash \varphi \rangle$  **have**  $(\Gamma \# \vdash n \varphi) = (n \leq \|\Gamma\|_\varphi)$

**by** (*simp add: relative-maximals-counting-deduction-lower-bound*)

**moreover have**  $(n \leq \|\Gamma\|_\varphi) = (\forall \Phi \in \mathcal{M} \Gamma \varphi. n \leq \text{length } (\Gamma \ominus \Phi))$

**proof** (*rule iffI*)

**assume**  $n \leq \|\Gamma\|_\varphi$

{

```

fix  $\Phi$ 
assume  $\Phi \in \mathcal{M} \Gamma \varphi$ 
hence  $n \leq \text{length} (\Gamma \ominus \Phi)$ 
  using  $\langle n \leq \|\Gamma\|_\varphi \rangle$  complement-relative-MaxSAT-intro by auto
}
thus  $\forall \Phi \in \mathcal{M} \Gamma \varphi. n \leq \text{length} (\Gamma \ominus \Phi)$  by blast
next
assume  $\forall \Phi \in \mathcal{M} \Gamma \varphi. n \leq \text{length} (\Gamma \ominus \Phi)$ 
with  $\langle \neg \vdash \varphi \rangle$  obtain  $\Phi$  where
   $\Phi \in \mathcal{M} \Gamma \varphi$ 
   $n \leq \text{length} (\Gamma \ominus \Phi)$ 
  using relative-maximals-existence
  by blast
thus  $n \leq \|\Gamma\|_\varphi$ 
  by (simp add: complement-relative-MaxSAT-intro)
qed
ultimately show ?thesis by metis
qed

```

**lemma** (*in consistent-classical-logic*) *counting-deduction-to-maxsat:*  
 $(\Gamma \# \vdash n \perp) = (\text{MaxSAT } \Gamma + n \leq \text{length } \Gamma)$   
**by** (*metis*  
*add.commute*  
*consistency*  
*length-MaxSAT-decomposition*  
*relative-maximals-counting-deduction-lower-bound*  
*nat-add-left-cancel-le*)

## Chapter 4

# Inequality Completeness For Probability Logic

### 4.1 Limited Counting Deduction Completeness

The reduction of counting deduction to MaxSAT allows us to first prove completeness for counting deduction, as maximal consistent sublists allow us to recover maximally consistent sets, which give rise to Dirac measures.

The completeness result first presented here, where all of the propositions on the left hand side are the same, will be extended later.

**lemma** (in *probability-logic*) *list-probability-upper-bound*:

$$(\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma) \leq \text{real } (\text{length } \Gamma)$$

**proof** (*induct*  $\Gamma$ )

case *Nil*

then show *?case* by *simp*

next

case (*Cons*  $\gamma$   $\Gamma$ )

moreover have  $\mathcal{P} \gamma \leq 1$  using *unity-upper-bound* by *blast*

ultimately have  $\mathcal{P} \gamma + (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma) \leq 1 + \text{real } (\text{length } \Gamma)$  by *linarith*

then show *?case* by *simp*

qed

**theorem** (in *classical-logic*) *dirac-limited-counting-deduction-completeness*:

$$(\forall \mathcal{P} \in \text{dirac-measures. } \text{real } n * \mathcal{P} \varphi \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) = \sim \Gamma \# \vdash n (\sim \varphi)$$

**proof** –

{

fix  $\mathcal{P} :: 'a \Rightarrow \text{real}$

assume  $\mathcal{P} \in \text{dirac-measures}$

from *this* interpret *probability-logic*  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$

unfolding *dirac-measures-def*

by *auto*

assume  $\sim \Gamma \# \vdash n (\sim \varphi)$

moreover have *replicate*  $n (\sim \varphi) = \sim (\text{replicate } n \varphi)$

by (*induct n, auto*)  
**ultimately have**  $\sim \Gamma \S \vdash \sim (\text{replicate } n \ \varphi)$   
 using *counting-deduction-to-measure-deduction by metis*  
**hence**  $(\sum \varphi \leftarrow (\text{replicate } n \ \varphi). \mathcal{P} \ \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \ \gamma)$   
 using *measure-deduction-soundness*  
 by *blast*  
**moreover have**  $(\sum \varphi \leftarrow (\text{replicate } n \ \varphi). \mathcal{P} \ \varphi) = \text{real } n * \mathcal{P} \ \varphi$   
 by (*induct n, simp, simp add: semiring-normalization-rules(3)*)  
**ultimately have**  $\text{real } n * \mathcal{P} \ \varphi \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \ \gamma)$   
 by *simp*  
 }  
**moreover**  
 {  
**assume**  $\neg \sim \Gamma \# \vdash n (\sim \varphi)$   
**have**  $\exists \mathcal{P} \in \text{dirac-measures. real } n * \mathcal{P} \ \varphi > (\sum \gamma \leftarrow \Gamma. \mathcal{P} \ \gamma)$   
**proof** –  
**have**  $\exists \Phi. \Phi \in \mathcal{M} (\sim \Gamma) (\sim \varphi)$   
 using  
    $\langle \neg \sim \Gamma \# \vdash n (\sim \varphi) \rangle$   
   *relative-maximals-existence*  
   *counting-deduction-tautology-weaken*  
 by *blast*  
**from this obtain**  $\Phi$  **where**  $\Phi$ :  
    $(\sim \Phi) \in \mathcal{M} (\sim \Gamma) (\sim \varphi)$   
    $\text{mset } \Phi \subseteq \# \text{ mset } \Gamma$   
**unfolding** *map-negation-def*  
 by (*metis*  
   (*mono-tags, lifting*)  
   *relative-maximals-def*  
   *mem-Collect-eq*  
   *mset-sub-map-list-exists*)  
**hence**  $\neg \vdash \varphi \rightarrow \sqcup \Phi$   
 using  
   *biconditional-weaken*  
   *list-deduction-def*  
   *map-negation-list-implication*  
   *set-deduction-base-theory*  
   *relative-maximals-def*  
 by *blast*  
**from this obtain**  $\Omega$  **where**  $\Omega$ :  $\text{MCS } \Omega \ \varphi \in \Omega \sqcup \Phi \notin \Omega$   
 by (*meson*  
   *insert-subset*  
   *formula-consistent-def*  
   *formula-maximal-consistency*  
   *formula-maximally-consistent-extension*  
   *formula-maximally-consistent-set-def-def*  
   *set-deduction-base-theory*  
   *set-deduction-reflection*  
   *set-deduction-theorem*)

```

let ?P = λ χ. if χ ∈ Ω then (1 :: real) else 0
from Ω have ?P ∈ dirac-measures
  using MCS-dirac-measure by blast
moreover
from this interpret probability-logic (λ φ. ⊢ φ) (→) ⊥ ?P
  unfolding dirac-measures-def
  by auto
have ∀ φ ∈ set Φ. ?P φ = 0
  using Φ(1) Ω(1) Ω(3) arbitrary-disjunction-exclusion-MCS by auto
with Φ(2) have (∑ γ ← Γ. ?P γ) = (∑ γ ← (Γ ⊖ Φ). ?P γ)
proof (induct Φ)
  case Nil
  then show ?case by simp
next
case (Cons φ Φ)
then show ?case
proof -
  obtain ω :: 'a where
    ω: ¬ mset Φ ⊆# mset Γ
      ∨ ω ∈ set Φ ∧ ω ∈ Ω
      ∨ (∑ γ ← Γ. ?P γ) = (∑ γ ← Γ ⊖ Φ. ?P γ)
  using Cons.hyps by fastforce
  have A:
    ∀ (f :: 'a ⇒ real) (Γ :: 'a list) Φ.
      ¬ mset Φ ⊆# mset Γ
      ∨ sum-list ((∑ φ ← Φ. f φ) # map f (Γ ⊖ Φ)) = (∑ γ ← Γ. f γ)
  using listSubtract-multisubset-list-summation by auto
  have B: ∀ rs. sum-list ((0 :: real) # rs) = sum-list rs
  by auto
  have C: ∀ r rs. (0 :: real) = r ∨ sum-list (r # rs) ≠ sum-list rs
  by simp
  have D: ∀ f. sum-list (sum-list (map f (φ # Φ)) # map f (Γ ⊖ (φ # Φ)))
    = (sum-list (map f Γ) :: real)
  using A Cons.prem(1) by blast
  have E: mset Φ ⊆# mset Γ
  using Cons.prem(1) subset-mset.dual-order.trans by force
  then have F: ∀ f. (0 :: real) = sum-list (map f Φ)
    ∨ sum-list (map f Γ) ≠ sum-list (map f (Γ ⊖ Φ))
  using C A by (metis (no-types))
  then have G: (∑ φ' ← (φ # Φ). ?P φ') = 0 ∨ ω ∈ Ω
  using E ω Cons.prem(2) by auto
  have H: ∀ Γ r :: real. r = (∑ γ ← Γ. ?P γ)
    ∨ ω ∈ set Φ
    ∨ r ≠ (∑ γ ← (φ # Γ). ?P γ)
  using Cons.prem(2) by auto
  have (1 :: real) ≠ 0 by linarith
moreover
{ assume ω ∉ set Φ
  then have ω ∉ Ω ∨ (∑ γ ← Γ. ?P γ) = (∑ γ ← Γ ⊖ (φ # Φ). ?P γ)

```



```

      using H F E D B ω by (metis (no-types) sum-list.Cons) }
    ultimately have ?thesis
      using G D B by (metis Cons.premis(2) list.set-intros(2))
    then show ?thesis
      by linarith
  qed
qed
hence (∑ γ←Γ. ?P γ) ≤ real (length (Γ ⊖ Φ))
  using list-probability-upper-bound
  by auto
  moreover
  have length (∼ Γ ⊖ ∼ Φ) < n
    by (metis not-le Φ(1) <¬ (∼ Γ) #+ n (∼ φ)
      relative-maximals-max-counting-deduction
      maximals-list-subtract-length-equiv)
  hence real (length (∼ Γ ⊖ ∼ Φ)) < real n
    by simp
  with Ω(2) have real (length (∼ Γ ⊖ ∼ Φ)) < real n * ?P φ
    by simp
  moreover
  have (∼ (Γ ⊖ Φ)) ≡ (∼ Γ ⊖ ∼ Φ)
    unfolding map-negation-def
    by (metis Φ(2) map-list-subtract-mset-equivalence)
  with perm-length have length (Γ ⊖ Φ) = length (∼ Γ ⊖ ∼ Φ)
    by (metis length-map local.map-negation-def)
  hence real (length (Γ ⊖ Φ)) = real (length (∼ Γ ⊖ ∼ Φ))
    by simp
  ultimately show ?thesis
    by force
  qed
}
ultimately show ?thesis by fastforce
qed

```

## 4.2 Measure Deduction Completeness

Since measure deduction may be reduced to counting deduction, we have measure deduction is complete.

**lemma** (in *classical-logic*) *dirac-measure-deduction-completeness*:

$$(\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) = \sim \Gamma \ \$\vdash \sim \Phi$$

**proof** –

```

{
  fix P :: 'a ⇒ real
  assume P ∈ dirac-measures
  from this interpret probability-logic (λ φ. ⊢ φ) (→) ⊥ P
  unfolding dirac-measures-def
  by auto
  assume ∼ Γ \$\vdash ∼ Φ

```

hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
 using *measure-deduction-soundness*  
 by *blast*  
 }  
 moreover  
 {  
 assume  $\neg \sim \Gamma \ \$\vdash \sim \Phi$   
 have  $\exists \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) > (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
 proof –  
 from  $\langle \neg \sim \Gamma \ \$\vdash \sim \Phi \rangle$  have  $\neg \sim (\sim \Phi) @ \sim \Gamma \ \#\vdash (\text{length } (\sim \Phi)) \perp$   
 using *measure-deduction-to-counting-deduction* by *blast*  
 moreover  
 have  $\sim (\sim \Phi) @ \sim \Gamma \ \#\vdash (\text{length } (\sim \Phi)) \perp = \sim (\sim \Phi) @ \sim \Gamma \ \#\vdash (\text{length } (\sim \Phi)) \perp$   
 }  
 by (*induct*  $\Phi$ , *auto*)  
 moreover have  $\vdash \sim \top \rightarrow \perp$   
 by (*simp add: negation-def*)  
 ultimately have  $\neg \sim (\sim \Phi @ \Gamma) \ \#\vdash (\text{length } \Phi) (\sim \top)$   
 using *counting-deduction-implication* by *fastforce*  
 from *this* obtain  $\mathcal{P}$  where  $\mathcal{P}$ :  
 $\mathcal{P} \in \text{dirac-measures}$   
 $\text{real } (\text{length } \Phi) * \mathcal{P} \top > (\sum \gamma \leftarrow (\sim \Phi @ \Gamma). \mathcal{P} \gamma)$   
 using *dirac-limited-counting-deduction-completeness*  
 by *fastforce*  
 from *this* interpret *probability-logic*  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$   
 unfolding *dirac-measures-def*  
 by *auto*  
 from  $\mathcal{P}(2)$  have  $\text{real } (\text{length } \Phi) > (\sum \gamma \leftarrow \sim \Phi. \mathcal{P} \gamma) + (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
 by (*simp add: probability-unity*)  
 moreover have  $(\sum \gamma \leftarrow \sim \Phi. \mathcal{P} \gamma) = \text{real } (\text{length } \Phi) - (\sum \gamma \leftarrow \Phi. \mathcal{P} \gamma)$   
 using *complementation*  
 by (*induct*  $\Phi$ , *auto*)  
 ultimately show *?thesis*  
 using  $\mathcal{P}(1)$  by *auto*  
 qed  
 }  
 ultimately show *?thesis* by *fastforce*  
 qed

**theorem** (*in classical-logic*) *measure-deduction-completeness*:  
 $(\forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) = \sim \Gamma \ \$\vdash \sim \Phi$   
 proof –  
 {  
 fix  $\mathcal{P} :: 'a \Rightarrow \text{real}$   
 assume  $\mathcal{P} \in \text{probabilities}$   
 from *this* interpret *probability-logic*  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$   
 unfolding *probabilities-def*  
 by *auto*  
 assume  $\sim \Gamma \ \$\vdash \sim \Phi$

hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
 using *measure-deduction-soundness*  
 by *blast*  
 }  
 thus *?thesis*  
 using *dirac-measures-subset dirac-measure-deduction-completeness*  
 by *fastforce*  
 qed

### 4.3 Counting Deduction Completeness

Leveraging our measure deduction completeness result, we may extend our limited counting deduction completeness theorem to full completeness.

**lemma** (in *classical-logic*) *measure-left-commute*:

$$(\Phi @ \Psi) \$\vdash \Xi = (\Psi @ \Phi) \$\vdash \Xi$$

**proof** –

$$\text{have } (\Phi @ \Psi) \preceq (\Psi @ \Phi) \quad (\Psi @ \Phi) \preceq (\Phi @ \Psi)$$

using *stronger-theory-reflexive stronger-theory-right-permutation perm-append-swap*

by *blast+*

thus *?thesis*

using *measure-stronger-theory-left-monotonic*

by *blast*

qed

**lemma** (in *classical-logic*) *stronger-theory-double-negation-right*:

$$\Phi \preceq \sim (\sim \Phi)$$

by (*induct*  $\Phi$ , *simp*, *simp add: double-negation negation-def stronger-theory-left-right-cons*)

**lemma** (in *classical-logic*) *stronger-theory-double-negation-left*:

$$\sim (\sim \Phi) \preceq \Phi$$

by (*induct*  $\Phi$ ,

*simp*,

*simp add: double-negation-converse negation-def stronger-theory-left-right-cons*)

**lemma** (in *classical-logic*) *counting-deduction-completeness*:

$$(\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) = (\sim \Gamma @ \Phi) \#\vdash (\text{length } \Phi) \perp$$

**proof** –

$$\text{have } (\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) \\ = \sim (\sim \Phi) @ \sim \Gamma \#\vdash (\text{length } (\sim \Phi)) \perp$$

using *dirac-measure-deduction-completeness measure-deduction-to-counting-deduction*

by *blast*

**also have** ... =  $\sim (\sim \Phi) @ \sim \Gamma \#\vdash (\text{length } \Phi) \perp$  by (*induct*  $\Phi$ , *auto*)

**also have** ... =  $\sim \Gamma @ \sim (\sim \Phi) \#\vdash (\text{length } \Phi) \perp$

by (*simp add: measure-left-commute counting-deduction-to-measure-deduction*)

**also have** ... =  $\sim \Gamma @ \Phi \#\vdash (\text{length } \Phi) \perp$

by (*meson measure-cancel*

*stronger-theory-to-measure-deduction*)

*measure-transitive*  
*counting-deduction-to-measure-deduction*  
*stronger-theory-double-negation-left*  
*stronger-theory-double-negation-right*  
**finally show** *?thesis by blast*  
**qed**

## 4.4 Collapse Theorem For Probability Logic

We now turn to proving the collapse theorem for probability logic. This states that any inequality holds for all finitely additive probability measures if and only if it holds for all Dirac measures.

**theorem** (in *classical-logic*) *weakly-additive-completeness-collapse*:  
 $(\forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
 $= (\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
**by** (*simp add: dirac-measure-deduction-completeness*  
*measure-deduction-completeness*)

The collapse theorem may be strengthened to include an arbitrary constant term  $c$ . This will be key to characterizing MaxSAT completeness in §4.5.

**lemma** (in *classical-logic*) *nat-dirac-probability*:  
 $\forall \mathcal{P} \in \text{dirac-measures}. \exists n :: \text{nat}. \text{real } n = (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi)$   
**proof** (*induct*  $\Phi$ )  
**case** *Nil*  
**then show** *?case by simp*  
**next**  
**case** (*Cons*  $\varphi \Phi$ )  
**{**  
**fix**  $\mathcal{P} :: 'a \Rightarrow \text{real}$   
**assume**  $\mathcal{P} \in \text{dirac-measures}$   
**from** *Cons this* **obtain**  $n$  **where**  $\text{real } n = (\sum \varphi' \leftarrow \Phi. \mathcal{P} \varphi')$  **by** *fastforce*  
**hence**  $\star: (\sum \varphi' \leftarrow \Phi. \mathcal{P} \varphi') = \text{real } n$  **by** *simp*  
**have**  $\exists n. \text{real } n = (\sum \varphi' \leftarrow (\varphi \# \Phi). \mathcal{P} \varphi')$   
**proof** (*cases*  $\mathcal{P} \varphi = 1$ )  
**case** *True*  
**then show** *?thesis*  
**by** (*simp add:  $\star$ ,metis of-nat-Suc*)  
**next**  
**case** *False*  
**hence**  $\mathcal{P} \varphi = 0$  **using**  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$  *dirac-measures-def* **by** *auto*  
**then show** *?thesis using  $\star$*   
**by** *simp*  
**qed**  
**}**  
**thus** *?case by blast*  
**qed**

**lemma** (in *classical-logic*) *dirac-ceiling*:  
 $\forall \mathcal{P} \in \text{dirac-measures.}$   
 $((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
 $= ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$

**proof** –  
{  
  **fix**  $\mathcal{P}$   
  **assume**  $\mathcal{P} \in \text{dirac-measures}$   
  **have**  $((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
   $= ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
  **proof** (*rule iffI*)  
  **assume** *assm*:  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
  **show**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
  **proof** (*rule ccontr*)  
  **assume**  $\neg (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
  **moreover**  
  **obtain**  $x :: \text{int}$   
  **and**  $y :: \text{int}$   
  **and**  $z :: \text{int}$   
  **where** *xyz*:  $x = (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi)$   
   $y = \lceil c \rceil$   
   $z = (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
  **using** *nat-dirac-probability*  
  **by** (*metis*  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$  *of-int-of-nat-eq*)  
  **ultimately have**  $x + y - 1 \geq z$  **by** *linarith*  
  **hence**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c > (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$  **using** *xyz* **by** *linarith*  
  **thus** *False* **using** *assm* **by** *simp*  
  **qed**  
  **next**  
  **assume**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
  **thus**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
  **by** *linarith*  
  **qed**  
}
**thus** *?thesis* **by** *blast*  
**qed**

**lemma** (in *probability-logic*) *probability-replicate-verum*:  
**fixes**  $n :: \text{nat}$   
**shows**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + n = (\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi)$   
**using** *probability-unity*  
**by** (*induct n, auto*)

**lemma** (in *classical-logic*) *dirac-collapse*:  
 $(\forall \mathcal{P} \in \text{probabilities. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
 $= (\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$

**proof**  
**assume**  $\forall \mathcal{P} \in \text{probabilities. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
**hence**  $\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$

```

    using dirac-measures-subset by fastforce
  thus  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
    using dirac-ceiling by blast
next
assume assm:  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
show  $\forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
proof (cases  $c \geq 0$ )
  case True
    from this obtain  $n :: \text{nat}$  where  $\text{real } n = \lceil c \rceil$ 
      by (metis (full-types)
        antisym-conv
        ceiling-le-zero
        ceiling-zero
        nat-0-iff
        nat-eq-iff2
        of-nat-nat)
    {
      fix  $\mathcal{P}$ 
      assume  $\mathcal{P} \in \text{dirac-measures}$ 
      from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
        unfolding dirac-measures-def
        by auto
      have  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
        using assm  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$  by blast
      hence  $(\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
        using  $\langle \text{real } n = \lceil c \rceil \rangle$ 
          probability-replicate-verum [where  $\Phi = \Phi$  and  $n = n$ ]
        by metis
    }
  hence  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
    by blast
  hence  $\dagger: \forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
    using weakly-additive-completeness-collapse by blast
  {
    fix  $\mathcal{P}$ 
    assume  $\mathcal{P} \in \text{probabilities}$ 
    from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
      unfolding probabilities-def
      by auto
    have  $(\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
      using  $\dagger$   $\langle \mathcal{P} \in \text{probabilities} \rangle$  by blast
    hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
      using  $\langle \text{real } n = \lceil c \rceil \rangle$ 
        probability-replicate-verum [where  $\Phi = \Phi$  and  $n = n$ ]
      by linarith
  }
}
then show ?thesis by blast

```

```

next
  case False
  hence  $\lceil c \rceil \leq 0$  by auto
  from this obtain  $n :: \text{nat}$  where  $\text{real } n = - \lceil c \rceil$ 
  by (metis neg-0-le-iff-le of-nat-nat)
  {
  fix  $\mathcal{P}$ 
  assume  $\mathcal{P} \in \text{dirac-measures}$ 
  from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
  unfolding dirac-measures-def
  by auto
  have  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
  using assm  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$  by blast
  hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$ 
  using  $\langle \text{real } n = - \lceil c \rceil \rangle$ 
  probability-replicate-verum [where  $\Phi = \Gamma$  and  $n = n$ ]
  by linarith
  }
  hence  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$ 
  by blast
  hence  $\ddagger: \forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$ 
  using weakly-additive-completeness-collapse by blast
  {
  fix  $\mathcal{P}$ 
  assume  $\mathcal{P} \in \text{probabilities}$ 
  from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
  unfolding probabilities-def
  by auto
  have  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$ 
  using  $\ddagger$   $\langle \mathcal{P} \in \text{probabilities} \rangle$  by blast
  hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
  using  $\langle \text{real } n = - \lceil c \rceil \rangle$ 
  probability-replicate-verum [where  $\Phi = \Gamma$  and  $n = n$ ]
  by linarith
  }
  then show ?thesis by blast
qed
qed

```

**lemma** (*in classical-logic*) *dirac-strict-floor*:

```

 $\forall \mathcal{P} \in \text{dirac-measures}. ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$ 
 $= ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$ 

```

**proof**

```

fix  $\mathcal{P} :: 'a \Rightarrow \text{real}$ 
let  $?P' = (\lambda \varphi. \lfloor \mathcal{P} \varphi \rfloor) :: 'a \Rightarrow \text{int}$ 
assume  $\mathcal{P} \in \text{dirac-measures}$ 

```

**hence**  $\forall \varphi. \mathcal{P} \varphi = ?\mathcal{P}' \varphi$   
**unfolding** *dirac-measures-def*  
**by** (*metis (mono-tags, lifting)*  
*mem-Collect-eq*  
*of-int-0*  
*of-int-1*  
*of-int-floor-cancel*)  
**hence**  $A: (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) = (\sum \varphi \leftarrow \Phi. ?\mathcal{P}' \varphi)$   
**by** (*induct*  $\Phi$ , *auto*)  
**have**  $B: (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma) = (\sum \gamma \leftarrow \Gamma. ?\mathcal{P}' \gamma)$   
**using**  $\langle \forall \varphi. \mathcal{P} \varphi = ?\mathcal{P}' \varphi \rangle$  **by** (*induct*  $\Gamma$ , *auto*)  
**have**  $((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
 $= ((\sum \varphi \leftarrow \Phi. ?\mathcal{P}' \varphi) + c < (\sum \gamma \leftarrow \Gamma. ?\mathcal{P}' \gamma))$   
**unfolding**  $A B$  **by** *auto*  
**also have**  $\dots = ((\sum \varphi \leftarrow \Phi. ?\mathcal{P}' \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. ?\mathcal{P}' \gamma))$   
**by** *linarith*  
**finally show**  $((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) =$   
 $((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
**using**  $A B$  **by** *linarith*  
**qed**

**lemma** (*in classical-logic*) *strict-dirac-collapse*:

$(\forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
 $= (\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$

**proof**

**assume**  $\forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
**hence**  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
**using** *dirac-measures-subset* **by** *blast*  
**thus**  $\forall \mathcal{P} \in \text{dirac-measures}. ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
**using** *dirac-strict-floor* **by** *blast*

**next**

**assume**  $\forall \mathcal{P} \in \text{dirac-measures}. ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
**moreover have**  $\lfloor c \rfloor + 1 = \lceil (\lfloor c \rfloor + 1) \rceil :: \text{real}$

**by** *simp*

**ultimately have**  $\star$ :

$\forall \mathcal{P} \in \text{probabilities}. ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
**using** *dirac-collapse [of*  $\Phi$   $\lfloor c \rfloor + 1$   $\Gamma]$

**by** *auto*

**show**  $\forall \mathcal{P} \in \text{probabilities}. ((\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$

**proof**

**fix**  $\mathcal{P} :: 'a \Rightarrow \text{real}$

**assume**  $\mathcal{P} \in \text{probabilities}$

**hence**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lfloor c \rfloor + 1 \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$

**using**  $\star$  **by** *auto*

**thus**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c < (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$

**by** *linarith*

**qed**

**qed**



## 4.5 MaxSAT Completeness For Probability Logic

It follows from the collapse theorem that any probability inequality tautology, include those with *constant terms*, may be reduced to a bounded MaxSAT problem. This is not only a key computational complexity result, but suggests a straightforward algorithm for *computing* probability identities.

**lemma** (in *classical-logic*) *relative-maximals-verum-extract*:

**assumes**  $\neg \vdash \varphi$

**shows**  $(| \text{replicate } n \top \ @ \ \Phi \ |_{\varphi}) = n + (| \Phi \ |_{\varphi})$

**proof** (*induct n*)

**case**  $0$

**then show** *?case by simp*

**next**

**case** (*Suc n*)

{

**fix**  $\Phi$

**obtain**  $\Sigma$  **where**  $\Sigma \in \mathcal{M} (\top \# \Phi) \varphi$

**using** *assms relative-maximals-existence* **by** *fastforce*

**hence**  $\top \in \text{set } \Sigma$

**by** (*metis (no-types, lifting)*)

*list.set-intros(1)*

*list-deduction-modus-ponens*

*list-deduction-weaken*

*relative-maximals-complement-equiv*

*relative-maximals-def*

*verum-tautology*

*mem-Collect-eq*)

**hence**  $\neg (\text{remove1 } \top \ \Sigma \vdash \varphi)$

**by** (*meson  $\langle \Sigma \in \mathcal{M} (\top \# \Phi) \varphi \rangle$* )

*list.set-intros(1)*

*axiom-k*

*list-deduction-modus-ponens*

*list-deduction-monotonic*

*list-deduction-weaken*

*relative-maximals-complement-equiv*

*set-remove1-subset*)

**moreover**

**have**  $\text{mset } \Sigma \subseteq \# \text{mset } (\top \# \Phi)$

**using**  $\langle \Sigma \in \mathcal{M} (\top \# \Phi) \varphi \rangle$  *relative-maximals-def* **by** *blast*

**hence**  $\text{mset } (\text{remove1 } \top \ \Sigma) \subseteq \# \text{mset } \Phi$

**using** *subset-eq-diff-conv* **by** *fastforce*

**ultimately have**  $(| \Phi \ |_{\varphi}) \geq \text{length } (\text{remove1 } \top \ \Sigma)$

**by** (*metis (no-types, lifting)*)

*relative-MaxSAT-intro*

*list-deduction-weaken*

*relative-maximals-def*

*relative-maximals-existence*  
*mem-Collect-eq*  
**hence**  $(|\Phi|_\varphi) + 1 \geq \text{length } \Sigma$   
**by** (*simp add: <math>\langle \top \in \text{set } \Sigma \rangle \text{length-remove1}</math>*)  
**moreover have**  $(|\Phi|_\varphi) < \text{length } \Sigma$   
**proof** (*rule ccontr*)  
**assume**  $\neg (|\Phi|_\varphi) < \text{length } \Sigma$   
**hence**  $(|\Phi|_\varphi) \geq \text{length } \Sigma$  **by** *linarith*  
**from this obtain**  $\Delta$  **where**  $\Delta \in \mathcal{M} \Phi \varphi$   $\text{length } \Delta \geq \text{length } \Sigma$   
**using** *assms relative-MaxSAT-intro relative-maximals-existence* **by** *fastforce*  
**hence**  $\neg (\top \# \Delta) \vdash \varphi$   
**using** *list-deduction-modus-ponens*  
*list-deduction-theorem*  
*list-deduction-weaken*  
*relative-maximals-def*  
*verum-tautology*  
**by** *blast*  
**moreover have**  $\text{mset } (\top \# \Delta) \subseteq\# \text{mset } (\top \# \Phi)$   
**using**  $\langle \Delta \in \mathcal{M} \Phi \varphi \rangle$  *relative-maximals-def* **by** *auto*  
**ultimately have**  $\text{length } \Sigma \geq \text{length } (\top \# \Delta)$   
**using**  $\langle \Sigma \in \mathcal{M} (\top \# \Phi) \varphi \rangle$  *relative-maximals-def* **by** *blast*  
**hence**  $\text{length } \Delta \geq \text{length } (\top \# \Delta)$   
**using**  $\langle \text{length } \Sigma \leq \text{length } \Delta \rangle$  *dual-order.trans* **by** *blast*  
**thus** *False* **by** *simp*  
**qed**  
**ultimately have**  $(|\top \# \Phi|_\varphi) = (1 + |\Phi|_\varphi)$   
**by** (*metis Suc-eq-plus1 Suc-le-eq <math>\langle \Sigma \in \mathcal{M} (\top \# \Phi) \varphi \rangle \text{add.commute le-antisym}</math>*  
*relative-MaxSAT-intro*)  
**}**  
**thus** *?case* **using** *Suc* **by** *simp*  
**qed**

**lemma** (*in classical-logic*) *complement-MaxSAT-completeness*:  
 $(\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) = (\text{length } \Phi \leq \|\sim \Gamma @ \Phi\|_\perp)$   
**proof** (*cases*  $\vdash \perp$ )  
**case** *True*  
**hence**  $\mathcal{M} (\sim \Gamma @ \Phi) \perp = \{\}$   
**using** *relative-maximals-existence* **by** *auto*  
**hence**  $\text{length } (\sim \Gamma @ \Phi) = \|\sim \Gamma @ \Phi\|_\perp$   
**unfolding** *complement-relative-MaxSAT-def* *relative-MaxSAT-def* **by** *presburger*  
**then show** *?thesis*  
**using** *True* *counting-deduction-completeness* *counting-deduction-tautology-weaken*  
**by** *auto*  
**next**  
**case** *False*  
**then show** *?thesis*  
**using** *counting-deduction-completeness* *relative-maximals-counting-deduction-lower-bound*

```

    by blast
qed

lemma (in classical-logic) relative-maximals-neg-verum-elim:
  (| replicate n ( $\sim \top$ ) @  $\Phi$  | $\varphi$ ) = (|  $\Phi$  | $\varphi$ )
proof (induct n)
  case 0
  then show ?case by simp
next
  case (Suc n)
  {
    fix  $\Phi$ 
    have (| ( $\sim \top$ ) #  $\Phi$  | $\varphi$ ) = (|  $\Phi$  | $\varphi$ )
    proof (cases  $\vdash \varphi$ )
      case True
      then show ?thesis
        unfolding relative-MaxSAT-def relative-maximals-def
        by (simp add: list-contradiction)
    next
      case False
      from this obtain  $\Sigma$  where  $\Sigma \in \mathcal{M} ((\sim \top) \# \Phi) \varphi$ 
      using relative-maximals-existence by fastforce
      have [ $\sim \top$ ]: $\vdash \varphi$ 
      by (metis modus-ponens
          Peirces-law
          pseudo-scotus
          list-contradiction
          list-contradiction-weaken
          negation-def
          verum-def)
      hence  $\sim \top \notin \text{set } \Sigma$ 
      by (meson  $\langle \Sigma \in \mathcal{M} (\sim \top \# \Phi) \varphi \rangle$ 
          list.set-intros(1)
          list-contradiction-base-theory
          list-contradiction-theorem
          list-contradiction-weaken
          relative-maximals-complement-equiv)
      hence remove1 ( $\sim \top$ )  $\Sigma = \Sigma$ 
      by (simp add: remove1-idem)
      moreover have mset  $\Sigma \subseteq\#$  mset (( $\sim \top$ ) #  $\Phi$ )
      using  $\langle \Sigma \in \mathcal{M} (\sim \top \# \Phi) \varphi \rangle$  relative-maximals-def by blast
      ultimately have mset  $\Sigma \subseteq\#$  mset  $\Phi$ 
      by (metis add-mset-add-single mset.simps(2) mset-remove1 subset-eq-diff-conv)
      moreover have  $\neg (\Sigma \vdash \varphi)$ 
      using  $\langle \Sigma \in \mathcal{M} (\sim \top \# \Phi) \varphi \rangle$  relative-maximals-def by blast
      ultimately have (|  $\Phi$  | $\varphi$ )  $\geq$  length  $\Sigma$ 
      by (metis (no-types, lifting)
          relative-MaxSAT-intro
          list-contradiction-weaken)
  }

```

```

      relative-maximals-def
      relative-maximals-existence
      mem-Collect-eq
hence (|  $\Phi$  | $\varphi$ )  $\geq$  (|  $(\sim \top) \# \Phi$  | $\varphi$ )
  using  $\langle \Sigma \in \mathcal{M} (\sim \top \# \Phi) \varphi \rangle$  relative-MaxSAT-intro by auto
moreover
have (|  $\Phi$  | $\varphi$ )  $\leq$  (|  $(\sim \top) \# \Phi$  | $\varphi$ )
proof –
  obtain  $\Delta$  where  $\Delta \in \mathcal{M} \Phi \varphi$ 
    using False relative-maximals-existence by blast
  hence
     $\neg \Delta : \vdash \varphi$ 
     $mset \Delta \subseteq \# mset ((\sim \top) \# \Phi)$ 
  unfolding relative-maximals-def
  by (simp,
    metis (mono-tags, lifting)
    Diff-eq-empty-iff-mset
    list-subtract.simps(2)
    list-subtract-mset-homomorphism
    relative-maximals-def
    mem-Collect-eq
    mset-zero-iff
    remove1.simps(1))
  hence  $length \Delta \leq length \Sigma$ 
    using  $\langle \Sigma \in \mathcal{M} (\sim \top \# \Phi) \varphi \rangle$  relative-maximals-def by blast
  thus ?thesis
    using  $\langle \Delta \in \mathcal{M} \Phi \varphi \rangle \langle \Sigma \in \mathcal{M} (\sim \top \# \Phi) \varphi \rangle$  relative-MaxSAT-intro by
auto
  qed
  ultimately show ?thesis
    using le-antisym by blast
  qed
}
thus ?case using Suc by simp
qed

```

**lemma** (in *classical-logic*) *dirac-MaxSAT-partial-completeness*:

$(\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)) = (\text{MaxSAT } (\sim \Gamma @ \Phi) \leq length \Gamma)$

**proof** –

```

{
  fix  $\mathcal{P} :: 'a \Rightarrow real$ 
  obtain  $\varrho :: 'a \text{ list} \Rightarrow 'a \text{ list} \Rightarrow 'a \Rightarrow real$  where
     $(\forall \Phi \Gamma. \varrho \Phi \Gamma \in \text{dirac-measures} \wedge \neg (\sum \varphi \leftarrow \Phi. (\varrho \Phi \Gamma) \varphi) \leq (\sum \gamma \leftarrow \Gamma. (\varrho \Phi \Gamma) \gamma))$ 
     $\vee length \Phi \leq \| \sim \Gamma @ \Phi \|_{\perp}$ 
     $\wedge (\forall \Phi \Gamma. length \Phi \leq (\| \sim \Gamma @ \Phi \|_{\perp})$ 
       $\longrightarrow (\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$ )
  using complement-MaxSAT-completeness by moura

```

**moreover have**  $\forall \Gamma \varphi n. \text{length } \Gamma - n \leq (\|\Gamma\|_\varphi) \vee (|\Gamma|_\varphi) - n \neq 0$   
**by** (*metis add-diff-cancel-right'*  
*cancel-ab-semigroup-add-class.diff-right-commute*  
*diff-is-0-eq length-MaxSAT-decomposition*)

**moreover have**  $\forall \Gamma \Phi n. \text{length } (\Gamma @ \Phi) - n \leq \text{length } \Gamma \vee \text{length } \Phi - n \neq 0$   
**by force**

**ultimately have**  
 $(\mathcal{P} \in \text{dirac-measures} \longrightarrow (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
 $\wedge (|\sim \Gamma @ \Phi|_\perp) \leq \text{length } (\sim \Gamma)$   
 $\vee \neg (|\sim \Gamma @ \Phi|_\perp) \leq \text{length } (\sim \Gamma)$   
 $\wedge (\exists \mathcal{P}. \mathcal{P} \in \text{dirac-measures} \wedge \neg (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$   
**by** (*metis (no-types) add-diff-cancel-left'*  
*add-diff-cancel-right'*  
*diff-is-0-eq length-append*  
*length-MaxSAT-decomposition*)

**}**  
**then show** *?thesis* **by auto**  
**qed**

**lemma** (*in consistent-classical-logic*) *dirac-inequality-elim*:  
**fixes**  $c :: \text{real}$   
**assumes**  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
**shows**  $(\text{MaxSAT } (\sim \Gamma @ \Phi) + c \leq \text{length } \Gamma)$   
**proof** (*cases*  $c \geq 0$ )  
**case True**  
**from this obtain**  $n :: \text{nat}$  **where**  $\text{real } n = \lceil c \rceil$   
**by** (*metis ceiling-mono ceiling-zero of-nat-nat*)  
**{**  
**fix**  $\mathcal{P}$   
**assume**  $\mathcal{P} \in \text{dirac-measures}$   
**from this interpret** *probability-logic*  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$   
**unfolding** *dirac-measures-def*  
**by auto**  
**have**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + n \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
**by** (*metis assms*  $\langle \mathcal{P} \in \text{dirac-measures} \rangle \langle \text{real } n = \lceil c \rceil \rangle$  *dirac-ceiling*)  
**hence**  $(\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
**using** *probability-replicate-verum* [**where**  $\Phi = \Phi$  **and**  $n = n$ ]  
**by metis**  
**}**  
**hence**  $(|\sim \Gamma @ \text{replicate } n \top @ \Phi|_\perp) \leq \text{length } \Gamma$   
**using** *dirac-MaxSAT-partial-completeness* **by blast**  
**moreover have**  $\text{mset } (\sim \Gamma @ \text{replicate } n \top @ \Phi) = \text{mset } (\text{replicate } n \top @ \sim \Gamma @ \Phi)$   
**by simp**  
**ultimately have**  $(|\text{replicate } n \top @ \sim \Gamma @ \Phi|_\perp) \leq \text{length } \Gamma$   
**unfolding** *relative-MaxSAT-def relative-maximals-def*  
**by metis**  
**hence**  $(|\sim \Gamma @ \Phi|_\perp) + \lceil c \rceil \leq \text{length } \Gamma$   
**using**  $\langle \text{real } n = \lceil c \rceil \rangle$  *consistency relative-maximals-verum-extract*

```

    by auto
  then show ?thesis by linarith
next
case False
hence  $\lceil c \rceil \leq 0$  by auto
from this obtain  $n :: \text{nat}$  where  $\text{real } n = - \lceil c \rceil$ 
  by (metis neg-0-le-iff-le of-nat-nat)
{
  fix  $\mathcal{P}$ 
  assume  $\mathcal{P} \in \text{dirac-measures}$ 
  from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
  unfolding dirac-measures-def
  by auto
  have  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
    using assms  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$  dirac-ceiling
  by blast
  hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma) + n$ 
    using  $\langle \text{real } n = - \lceil c \rceil \rangle$  by linarith
  hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$ 
    using probability-replicate-verum [where  $\Phi = \Gamma$  and  $n = n$ ]
    by metis
}
hence  $(|\sim (\text{replicate } n \top @ \Gamma) @ \Phi |_{\perp}) \leq \text{length } (\text{replicate } n \top @ \Gamma)$ 
  using dirac-MaxSAT-partial-completeness [where  $\Phi = \Phi$  and  $\Gamma = \text{replicate } n \top$ 
@  $\Gamma$ ]
  by metis
hence  $(|\sim \Gamma @ \Phi |_{\perp}) \leq n + \text{length } \Gamma$ 
  by (simp add: relative-maximals-neg-verum-elim)
then show ?thesis using  $\langle \text{real } n = - \lceil c \rceil \rangle$  by linarith
qed

lemma (in classical-logic) dirac-inequality-intro:
  fixes  $c :: \text{real}$ 
  assumes  $\text{MaxSAT } (\sim \Gamma @ \Phi) + c \leq \text{length } \Gamma$ 
  shows  $\forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
proof (cases  $\vdash \perp$ )
  assume  $\vdash \perp$ 
  {
    fix  $\mathcal{P}$ 
    assume  $\mathcal{P} \in \text{dirac-measures}$ 
    from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
    unfolding dirac-measures-def
    by auto
    have False
      using  $\langle \vdash \perp \rangle$  consistency by blast
  }
  then show ?thesis by blast
next
assume  $\neg \vdash \perp$ 

```

**then show** *?thesis*  
**proof** (cases  $c \geq 0$ )  
    **assume**  $c \geq 0$   
    **from this obtain**  $n :: \text{nat}$  **where**  $\text{real } n = \lceil c \rceil$   
    by (metis ceiling-mono ceiling-zero of-nat-nat)  
    **hence**  $n + (|\sim \Gamma @ \Phi |_{\perp}) \leq \text{length } \Gamma$   
    **using** *assms* **by** *linarith*  
    **hence**  $(|\text{replicate } n \top @ \sim \Gamma @ \Phi |_{\perp}) \leq \text{length } \Gamma$   
    by (*simp add: <math>\langle \neg \vdash \perp \rangle relative-maximals-verum-extract*)  
    **moreover have**  $\text{mset } (\text{replicate } n \top @ \sim \Gamma @ \Phi) = \text{mset } (\sim \Gamma @ \text{replicate } n \top @ \Phi)$   
    by *simp*  
    **ultimately have**  $(|\sim \Gamma @ \text{replicate } n \top @ \Phi |_{\perp}) \leq \text{length } \Gamma$   
    **unfolding** *relative-MaxSAT-def* *relative-maximals-def*  
    **by** *metis*  
    **hence**  $\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
    **using** *dirac-MaxSAT-partial-completeness* **by** *blast*  
    {  
    fix  $\mathcal{P}$   
    **assume**  $\mathcal{P} \in \text{dirac-measures}$   
    **from this interpret** *probability-logic*  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$   
    **unfolding** *dirac-measures-def*  
    **by** *auto*  
    **have**  $(\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
    **using**  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$   
     $\langle \forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow (\text{replicate } n \top) @ \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma) \rangle$   
    **by** *blast*  
    **hence**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + n \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
    **by** (*simp add: probability-replicate-verum*)  
    **hence**  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$   
    **using**  $\langle \text{real } n = \text{real-of-int } \lceil c \rceil \rangle$  **by** *linarith*  
    }  
    **then show** *?thesis* **by** *blast*  
**next**  
    **assume**  $\neg (c \geq 0)$   
    **hence**  $\lceil c \rceil \leq 0$  **by** *auto*  
    **from this obtain**  $n :: \text{nat}$  **where**  $\text{real } n = -\lceil c \rceil$   
    by (metis neg-0-le-iff-le of-nat-nat)  
    **hence**  $(|\sim \Gamma @ \Phi |_{\perp}) \leq n + \text{length } \Gamma$   
    **using** *assms* **by** *linarith*  
    **hence**  $(|\sim (\text{replicate } n \top @ \Gamma) @ \Phi |_{\perp}) \leq \text{length } (\text{replicate } n \top @ \Gamma)$   
    by (*simp add: relative-maximals-neg-verum-elim*)  
    **hence**  $\forall \mathcal{P} \in \text{dirac-measures. } (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$   
    **using** *dirac-MaxSAT-partial-completeness* **by** *blast*  
    {  
    fix  $\mathcal{P}$

```

assume  $\mathcal{P} \in \text{dirac-measures}$ 
from this interpret probability-logic  $(\lambda \varphi. \vdash \varphi) (\rightarrow) \perp \mathcal{P}$ 
  unfolding dirac-measures-def
  by auto
have  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma)$ 
  using  $\langle \mathcal{P} \in \text{dirac-measures} \rangle$ 
     $\langle \forall \mathcal{P} \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) \leq (\sum \gamma \leftarrow (\text{replicate } n \top) @ \Gamma. \mathcal{P} \gamma) \rangle$ 
  by blast
hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + \lceil c \rceil \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
  using  $\langle \text{real } n = - \lceil c \rceil \rangle$  probability-replicate-verum by auto
hence  $(\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma)$ 
  by linarith
}
then show ?thesis by blast
qed
qed

lemma (in consistent-classical-logic) dirac-inequality-equiv:
   $(\forall \delta \in \text{dirac-measures}. (\sum \varphi \leftarrow \Phi. \delta \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \delta \gamma))$ 
   $= (\text{MaxSAT } (\sim \Gamma @ \Phi) + (c :: \text{real}) \leq \text{length } \Gamma)$ 
using dirac-inequality-elim dirac-inequality-intro consistency by auto

theorem (in consistent-classical-logic) probability-inequality-equiv:
   $(\forall \mathcal{P} \in \text{probabilities}. (\sum \varphi \leftarrow \Phi. \mathcal{P} \varphi) + c \leq (\sum \gamma \leftarrow \Gamma. \mathcal{P} \gamma))$ 
   $= (\text{MaxSAT } (\sim \Gamma @ \Phi) + (c :: \text{real}) \leq \text{length } \Gamma)$ 
unfolding dirac-collapse
using dirac-inequality-equiv dirac-ceiling by blast

no-notation first-component  $\langle \mathfrak{A} \rangle$ 
no-notation second-component  $\langle \mathfrak{B} \rangle$ 
no-notation merge-witness  $\langle \mathfrak{J} \rangle$ 
no-notation X-witness  $\langle \mathfrak{X} \rangle$ 
no-notation X-component  $\langle \mathfrak{X}_\bullet \rangle$ 
no-notation Y-witness  $\langle \mathfrak{Y} \rangle$ 
no-notation Y-component  $\langle \mathfrak{Y}_\bullet \rangle$ 
no-notation submerge-witness  $\langle \mathfrak{E} \rangle$ 
no-notation recover-witness-A  $\langle \mathfrak{P} \rangle$ 
no-notation recover-complement-A  $\langle \mathfrak{P}^C \rangle$ 
no-notation recover-witness-B  $\langle \mathfrak{Q} \rangle$ 
no-notation relative-maximals  $\langle \mathfrak{M} \rangle$ 
no-notation relative-MaxSAT  $\langle | - | \rightarrow [45] \rangle$ 
no-notation complement-relative-MaxSAT  $\langle \| - \| \rightarrow [45] \rangle$ 
no-notation MaxSAT-optimal-pre-witness  $\langle \mathfrak{W} \rangle$ 
no-notation MaxSAT-optimal-witness  $\langle \mathfrak{W} \rangle$ 
no-notation disjunction-MaxSAT-optimal-witness  $\langle \mathfrak{W}_{\sqcup} \rangle$ 
no-notation implication-MaxSAT-optimal-witness  $\langle \mathfrak{W}_{\rightarrow} \rangle$ 
no-notation MaxSAT-witness  $\langle \mathfrak{W} \rangle$ 

```



**unbundle** *funcset-syntax*

**end**

# Bibliography

- [1] M. R. Garey, D. S. Johnson, and L. Stockmeyer. Some simplified NP-complete graph problems. *Theoretical Computer Science*, 1(3):237–267, Feb. 1976.