

Probabilistic Hierarchy

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1 Bisimilarity

definition *bisimilar* where

$$\textit{bisimilar } Q \textit{ } s1 \textit{ } s2 \textit{ } x \textit{ } y \equiv (\exists R. R \textit{ } x \textit{ } y \wedge (\forall x \textit{ } y. R \textit{ } x \textit{ } y \longrightarrow Q \textit{ } R \textit{ } (s1 \textit{ } x) \textit{ } (s2 \textit{ } y)))$$

abbreviation *bisimilar-mc* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-pmf } R)$

abbreviation *bisimilar-dlts* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-fun } (=) \textit{ (rel-option } R))$

abbreviation *bisimilar-lts* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-bset } (\textit{rel-prod } (=) \textit{ } R))$

abbreviation *bisimilar-react* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-fun } (=) \textit{ (rel-option } (\textit{rel-pmf } R)))$

abbreviation *bisimilar-lmc* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-prod } (=) \textit{ (rel-pmf } R))$

abbreviation *bisimilar-lmdp* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-prod } (=) \textit{ (rel-nebset } (\textit{rel-pmf } R)))$

abbreviation *bisimilar-gen* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-option } (\textit{rel-pmf } (\textit{rel-prod } (=) \textit{ } R)))$

abbreviation $\text{bisimilar-str} \equiv \text{bisimilar } (\lambda R. \text{rel-sum } (\text{rel-pmf } R) (\text{rel-option } (\text{rel-prod } (=) R)))$

abbreviation $\text{bisimilar-alt} \equiv \text{bisimilar } (\lambda R. \text{rel-sum } (\text{rel-pmf } R) (\text{rel-bset } (\text{rel-prod } (=) R)))$

abbreviation $\text{bisimilar-sseg} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-prod } (=) (\text{rel-pmf } R)))$

abbreviation $\text{bisimilar-seg} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-pmf } (\text{rel-prod } (=) R)))$

abbreviation $\text{bisimilar-bun} \equiv \text{bisimilar } (\lambda R. \text{rel-pmf } (\text{rel-bset } (\text{rel-prod } (=) R)))$

abbreviation $\text{bisimilar-pz} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-pmf } (\text{rel-bset } (\text{rel-prod } (=) R))))$

abbreviation $\text{bisimilar-mg} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-pmf } (\text{rel-bset } (\text{rel-sum } (\text{rel-prod } (=) R) R))))$

2 Systems

codatatype $\text{mc} = \text{MC } \text{mc } \text{pmf}$

codatatype $'a \text{ dlts} = \text{DLTS } 'a \Rightarrow 'a \text{ dlts } \text{option}$

codatatype $('a, 'k) \text{ lts} = \text{LTS } ('a \times ('a, 'k) \text{ lts}) \text{ set}['k]$

codatatype $'a \text{ react} = \text{React } 'a \Rightarrow 'a \text{ react } \text{pmf } \text{option}$

codatatype $'a \text{ lmc} = \text{LMC } 'a \times 'a \text{ lmc } \text{pmf}$

codatatype $('a, 'k) \text{ lmdp} = \text{LMDP } 'a \times ('a, 'k) \text{ lmdp } \text{pmf } \text{set}!['k]$

codatatype $'a \text{ gen} = \text{Gen } ('a \times 'a \text{ gen}) \text{pmf } \text{option}$

codatatype $'a \text{ str} = \text{Str } 'a \text{ str } \text{pmf} + ('a \times 'a \text{ str}) \text{option}$

codatatype $('a, 'k) \text{ alt} = \text{Alt } ('a, 'k) \text{ alt } \text{pmf} + ('a \times ('a, 'k) \text{ alt}) \text{set}['k]$

codatatype $('a, 'k) \text{ sseg} = \text{SSeg } ('a \times ('a, 'k) \text{ sseg } \text{pmf}) \text{set}['k]$

codatatype $('a, 'k) \text{ seg} = \text{Seg } ('a \times ('a, 'k) \text{ seg}) \text{pmf } \text{set}['k]$

codatatype $('a, 'k) \text{ bun} = \text{Bun } (('a \times ('a, 'k) \text{ bun}) \text{set}['k]) \text{pmf}$

codatatype $('a, 'k1, 'k2) \text{ pz} = \text{PZ } (('a \times ('a, 'k1, 'k2) \text{ pz}) \text{set}['k1]) \text{pmf } \text{set}['k2]$

codatatype $('a, 'k1, 'k2) \text{ mg} = \text{MG } (('a \times ('a, 'k1, 'k2) \text{ mg} + ('a, 'k1, 'k2) \text{ mg}) \text{set}['k1]) \text{pmf } \text{set}['k2]$

3 Unfolds

primcorec $\text{unfold-mc} :: ('a \Rightarrow 'a \text{ pmf}) \Rightarrow 'a \Rightarrow \text{mc } \mathbf{where}$
 $\text{unfold-mc } s \ x = \text{MC } (\text{map-pmf } (\text{unfold-mc } s) (s \ x))$

primcorec $\text{unfold-dlts} :: ('a \Rightarrow 'b \Rightarrow 'a \text{ option}) \Rightarrow 'a \Rightarrow 'b \text{ dlts } \mathbf{where}$
 $\text{unfold-dlts } s \ x = \text{DLTS } (\text{map-option } (\text{unfold-dlts } s) o \ s \ x)$

primcorec $\text{unfold-lts} :: ('a \Rightarrow ('b \times 'a) \text{ set}['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ lts } \mathbf{where}$
 $\text{unfold-lts } s \ x = \text{LTS } (\text{map-bset } (\text{map-prod } \text{id } (\text{unfold-lts } s)) (s \ x))$

primcorec $\text{unfold-react} :: ('a \Rightarrow 'b \Rightarrow 'a \text{ pmf } \text{option}) \Rightarrow 'a \Rightarrow 'b \text{ react } \mathbf{where}$
 $\text{unfold-react } s \ x = \text{React } (\text{map-option } (\text{map-pmf } (\text{unfold-react } s)) o \ s \ x)$

primcorec $\text{unfold-lmc} :: ('a \Rightarrow 'b \times 'a \text{ pmf}) \Rightarrow 'a \Rightarrow 'b \text{ lmc } \mathbf{where}$
 $\text{unfold-lmc } s \ x = \text{LMC } (\text{map-prod } \text{id } (\text{map-pmf } (\text{unfold-lmc } s)) (s \ x))$

primcorec $\text{unfold-lmdp} :: ('a \Rightarrow 'b \times 'a \text{ pmf } \text{set}!['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ lmdp } \mathbf{where}$

unfold-lmdp $s x = LMDP (map-prod id (map-nebset (map-pmf (unfold-lmdp s))) (s x))$

primcorec *unfold-gen* $:: ('a \Rightarrow (('b \times 'a) pmf) option) \Rightarrow 'a \Rightarrow 'b \text{ gen where}$
unfold-gen $s x = Gen (map-option (map-pmf (map-prod id (unfold-gen s))) (s x))$

primcorec *unfold-str* $:: ('a \Rightarrow 'a pmf + ('b \times 'a) option) \Rightarrow 'a \Rightarrow 'b \text{ str where}$
unfold-str $s x = Str (map-sum (map-pmf (unfold-str s)) (map-option (map-prod id (unfold-str s)))) (s x)$

primcorec *unfold-alt* $:: ('a \Rightarrow 'a pmf + ('b \times 'a) set['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ alt where}$
unfold-alt $s x = Alt (map-sum (map-pmf (unfold-alt s)) (map-bset (map-prod id (unfold-alt s)))) (s x)$

primcorec *unfold-sseg* $:: ('a \Rightarrow ('b \times 'a pmf) set['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ sseg where}$
unfold-sseg $s x = SSeg (map-bset (map-prod id (map-pmf (unfold-sseg s)))) (s x)$

primcorec *unfold-seg* $:: ('a \Rightarrow (('b \times 'a) pmf) set['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ seg where}$
unfold-seg $s x = Seg (map-bset (map-pmf (map-prod id (unfold-seg s)))) (s x)$

primcorec *unfold-bun* $:: ('a \Rightarrow (('b \times 'a) set['k]) pmf) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ bun where}$
unfold-bun $s x = Bun (map-pmf (map-bset (map-prod id (unfold-bun s)))) (s x)$

primcorec *unfold-pz* $:: ('a \Rightarrow (('b \times 'a) set['k1]) pmf set['k2]) \Rightarrow 'a \Rightarrow ('b, 'k1, 'k2) \text{ pz where}$
unfold-pz $s x = PZ (map-bset (map-pmf (map-bset (map-prod id (unfold-pz s)))) (s x))$

primcorec *unfold-mg* $:: ('a \Rightarrow (('b \times 'a + 'a) set['k1]) pmf set['k2]) \Rightarrow 'a \Rightarrow ('b, 'k1, 'k2) \text{ mg where}$
unfold-mg $s x = MG (map-bset (map-pmf (map-bset (map-sum (map-prod id (unfold-mg s)) (unfold-mg s)))) (s x))$

4 Embeddings

abbreviation (*input*) *react-of-dlts-emb* $dlts \equiv map-option \text{ return-pmf } o \text{ dlts}$

abbreviation (*input*) *lts-of-dlts-emb* $\equiv bgraph$

abbreviation (*input*) *sseg-of-react-emb* $\equiv bgraph$

abbreviation (*input*) *gen-of-lmc-emb* $\equiv Some \text{ o case-prod } (map-pmf \text{ o } Pair)$

abbreviation (*input*) *lmdp-of-lmc-emb* $\equiv map-prod \text{ id } \text{ nebsingleton}$

abbreviation (*input*) *sseg-of-lmdp-emb* $\equiv (\lambda(a, X). map-bset (Pair a) (bset-of-nebset X))$

abbreviation (*input*) *sseg-of-lts-emb* $\equiv map-bset (map-prod \text{ id } \text{ return-pmf})$

abbreviation (*input*) *ssegopt-of-alt-emb* $\equiv case-sum$

$(map-bset (Pair None) \text{ o } \text{ bsingleton})$

$(map-bset (map-prod \text{ Some } \text{ return-pmf}))$

abbreviation (*input*) *bunopt-of-alt-emb* $\equiv case-sum$

$(\text{map-pmf } (\text{bsingleton } o \text{ Pair None}))$
 $(\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id)) o \text{ return-pmf})$
abbreviation $(\text{input}) \text{ segopt-of-seg-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-prod } \text{Some } id))$
abbreviation $(\text{input}) \text{ ssegopt-of-sseg-emb} \equiv \text{map-bset } (\text{map-prod } \text{Some } id)$
abbreviation $(\text{input}) \text{ bunopt-of-bun-emb} \equiv \text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id))$
abbreviation $(\text{input}) \text{ pzopt-of-pz-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id)))$
abbreviation $(\text{input}) \text{ seg-of-sseg-emb} \equiv \text{map-bset } (\text{case-prod } (\text{map-pmf } o \text{ Pair}))$
abbreviation $(\text{input}) \text{ pz-of-seg-emb} \equiv \text{map-bset } (\text{map-pmf } \text{bsingleton})$
abbreviation $(\text{input}) \text{ pz-of-bun-emb} \equiv \text{bsingleton}$
abbreviation $(\text{input}) \text{ seg-of-gen-emb} \equiv \text{bset-of-option}$
abbreviation $(\text{input}) \text{ bun-of-lts-emb} \equiv \text{return-pmf}$
abbreviation $(\text{input}) \text{ bun-of-gen-emb} \equiv \text{case-option } (\text{return-pmf } \text{bempty}) (\text{map-pmf } \text{bsingleton})$
abbreviation $(\text{input}) \text{ str-of-mc-emb} \equiv \text{Inl}$
abbreviation $(\text{input}) \text{ alt-of-str-emb} \equiv \text{map-sum } id \text{ bset-of-option}$
abbreviation $(\text{input}) \text{ pzopt-of-mg-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{case-sum } (\text{map-prod } \text{Some } id) (\text{Pair } \text{None}))))$
abbreviation $(\text{input}) \text{ mg-of-pzopt-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } (\lambda(a, s). \text{case-option } (\text{Inr } s) (\lambda a. (\text{Inl } (a, s))) a))))$

Obsolete edges (susumed by transitive ones)

abbreviation $(\text{input}) \text{ mg-of-pz-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inl}))$
abbreviation $(\text{input}) \text{ mg-of-alt1-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inr } o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inl } o \text{ bsingleton}) o \text{ return-pmf}))$
abbreviation $(\text{input}) \text{ mg-of-alt2-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inr } o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inl}) o \text{ return-pmf}) o \text{ bsingleton})$
abbreviation $(\text{input}) \text{ pz-of-alt1-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{Pair } \text{None}) o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id) o \text{ bsingleton}) o \text{ return-pmf}))$
abbreviation $(\text{input}) \text{ pz-of-alt2-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{Pair } \text{None}) o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id)) o \text{ return-pmf}) o \text{ bsingleton})$

definition $\text{react-of-dlts} :: 'a \text{ dlts} \Rightarrow 'a \text{ react}$ **where**
 $[\text{simp}]: \text{react-of-dlts} = \text{unfold-react } (\text{react-of-dlts-emb } o \text{ un-DLTS})$

definition $\text{lts-of-dlts} :: 'a \text{ dlts} \Rightarrow ('a, 'a \text{ set}) \text{ lts}$ **where**
 $[\text{simp}]: \text{lts-of-dlts} = \text{unfold-lts } (\text{lts-of-dlts-emb } o \text{ un-DLTS})$

definition $\text{sseg-of-react} :: 'a \text{ react} \Rightarrow ('a, 'a \text{ set}) \text{ sseg}$ **where**
 $[\text{simp}]: \text{sseg-of-react} = \text{unfold-sseg } (\text{sseg-of-react-emb } o \text{ un-React})$

definition $\text{lmdp-of-lmc} :: 'a \text{ lmc} \Rightarrow ('a, 'k) \text{ lmdp}$ **where**

[simp]: $lmdp\text{-of}\text{-}lmc = \text{unfold}\text{-}lmdp (\text{lmdp}\text{-of}\text{-}lmc\text{-emb } o \text{ un-LMC})$

definition $gen\text{-of}\text{-}lmc :: 'a \text{ lmc} \Rightarrow 'a \text{ gen}$ **where**

[simp]: $gen\text{-of}\text{-}lmc = \text{unfold}\text{-}gen (\text{gen}\text{-of}\text{-}lmc\text{-emb } o \text{ un-LMC})$

definition $sseg\text{-of}\text{-}lmdp :: ('a, 'k) \text{ lmdp} \Rightarrow ('a, 'k) \text{ sseg}$ **where**

[simp]: $sseg\text{-of}\text{-}lmdp = \text{unfold}\text{-}sseg (\text{sseg}\text{-of}\text{-}lmdp\text{-emb } o \text{ un-LMDP})$

definition $sseg\text{-of}\text{-}lts :: ('a, 'k) \text{ lts} \Rightarrow ('a, 'k) \text{ sseg}$ **where**

[simp]: $sseg\text{-of}\text{-}lts = \text{unfold}\text{-}sseg (\text{sseg}\text{-of}\text{-}lts\text{-emb } o \text{ un-LTS})$

definition $ssegopt\text{-of}\text{-}alt :: ('a, 'k) \text{ alt} \Rightarrow ('a \text{ option}, 'k) \text{ sseg}$ **where**

[simp]: $ssegopt\text{-of}\text{-}alt = \text{unfold}\text{-}sseg (\text{ssegopt}\text{-of}\text{-}alt\text{-emb } o \text{ un-Alt})$

definition $bunopt\text{-of}\text{-}alt :: ('a, 'k) \text{ alt} \Rightarrow ('a \text{ option}, 'k) \text{ bun}$ **where**

[simp]: $bunopt\text{-of}\text{-}alt = \text{unfold}\text{-}bun (\text{bunopt}\text{-of}\text{-}alt\text{-emb } o \text{ un-Alt})$

definition $seg\text{-of}\text{-}sseg :: ('a, 'k) \text{ sseg} \Rightarrow ('a, 'k) \text{ seg}$ **where**

[simp]: $seg\text{-of}\text{-}sseg = \text{unfold}\text{-}seg (\text{seg}\text{-of}\text{-}sseg\text{-emb } o \text{ un-SSeg})$

definition $seg\text{-of}\text{-}gen :: 'a \text{ gen} \Rightarrow ('a, 'k) \text{ seg}$ **where**

[simp]: $seg\text{-of}\text{-}gen = \text{unfold}\text{-}seg (\text{seg}\text{-of}\text{-}gen\text{-emb } o \text{ un-Gen})$

definition $bun\text{-of}\text{-}lts :: ('a, 'k) \text{ lts} \Rightarrow ('a, 'k) \text{ bun}$ **where**

[simp]: $bun\text{-of}\text{-}lts = \text{unfold}\text{-}bun (\text{bun}\text{-of}\text{-}lts\text{-emb } o \text{ un-LTS})$

definition $bun\text{-of}\text{-}gen :: 'a \text{ gen} \Rightarrow ('a, 'k) \text{ bun}$ **where**

[simp]: $bun\text{-of}\text{-}gen = \text{unfold}\text{-}bun (\text{bun}\text{-of}\text{-}gen\text{-emb } o \text{ un-Gen})$

definition $pz\text{-of}\text{-}seg :: ('a, 'k) \text{ seg} \Rightarrow ('a, 'k1, 'k) \text{ pz}$ **where**

[simp]: $pz\text{-of}\text{-}seg = \text{unfold}\text{-}pz (\text{pz}\text{-of}\text{-}seg\text{-emb } o \text{ un-Seg})$

definition $pz\text{-of}\text{-}bun :: ('a, 'k) \text{ bun} \Rightarrow ('a, 'k, 'k1) \text{ pz}$ **where**

[simp]: $pz\text{-of}\text{-}bun = \text{unfold}\text{-}pz (\text{pz}\text{-of}\text{-}bun\text{-emb } o \text{ un-Bun})$

definition $mg\text{-of}\text{-}pz :: ('a, 'k1, 'k2) \text{ pz} \Rightarrow ('a, 'k1, 'k2) \text{ mg}$ **where**

[simp]: $mg\text{-of}\text{-}pz = \text{unfold}\text{-}mg (\text{mg}\text{-of}\text{-}pz\text{-emb } o \text{ un-PZ})$

definition $str\text{-of}\text{-}mc :: mc \Rightarrow 'a \text{ str}$ **where**

[simp]: $str\text{-of}\text{-}mc = \text{unfold}\text{-}str (\text{str}\text{-of}\text{-}mc\text{-emb } o \text{ un-MC})$

definition $alt\text{-of}\text{-}str :: 'a \text{ str} \Rightarrow ('a, 'k) \text{ alt}$ **where**

[simp]: $alt\text{-of}\text{-}str = \text{unfold}\text{-}alt (\text{alt}\text{-of}\text{-}str\text{-emb } o \text{ un-Str})$

definition $ssegopt\text{-of}\text{-}sseg :: ('a, 'k) \text{ sseg} \Rightarrow ('a \text{ option}, 'k) \text{ sseg}$ **where**

[simp]: $ssegopt\text{-of}\text{-}sseg = \text{unfold}\text{-}sseg (\text{ssegopt}\text{-of}\text{-}sseg\text{-emb } o \text{ un-SSeg})$

definition $segopt\text{-of}\text{-}seg :: ('a, 'k) \text{ seg} \Rightarrow ('a \text{ option}, 'k) \text{ seg}$ **where**

[simp]: $segopt\text{-of}\text{-}seg = \text{unfold}\text{-}seg (\text{segopt}\text{-of}\text{-}seg\text{-emb } o \text{ un-Seg})$

definition $\text{bunopt-of-bun} :: ('a, 'k) \text{ bun} \Rightarrow ('a \text{ option}, 'k) \text{ bun}$ **where**

[simp]: $\text{bunopt-of-bun} = \text{unfold-bun} (\text{bunopt-of-bun-emb} \circ \text{un-Bun})$

definition $\text{pzopt-of-pz} :: ('a, 'k1, 'k2) \text{ pz} \Rightarrow ('a \text{ option}, 'k1, 'k2) \text{ pz}$ **where**

[simp]: $\text{pzopt-of-pz} = \text{unfold-pz} (\text{pzopt-of-pz-emb} \circ \text{un-PZ})$

definition $\text{pzopt-of-mg} :: ('a, 'k1, 'k2) \text{ mg} \Rightarrow ('a \text{ option}, 'k1, 'k2) \text{ pz}$ **where**

[simp]: $\text{pzopt-of-mg} = \text{unfold-pz} (\text{pzopt-of-mg-emb} \circ \text{un-MG})$

definition $\text{mg-of-pzopt} :: ('a \text{ option}, 'k1, 'k2) \text{ pz} \Rightarrow ('a, 'k1, 'k2) \text{ mg}$ **where**

[simp]: $\text{mg-of-pzopt} = \text{unfold-mg} (\text{mg-of-pzopt-emb} \circ \text{un-PZ})$

definition $\text{mg-of-alt1} :: ('a, 'k) \text{ alt} \Rightarrow ('a, 'k1, 'k) \text{ mg}$ **where**

[simp]: $\text{mg-of-alt1} = \text{unfold-mg} (\text{mg-of-alt1-emb} \circ \text{un-Alt})$

definition $\text{mg-of-alt2} :: ('a, 'k) \text{ alt} \Rightarrow ('a, 'k, 'k1) \text{ mg}$ **where**

[simp]: $\text{mg-of-alt2} = \text{unfold-mg} (\text{mg-of-alt2-emb} \circ \text{un-Alt})$

definition $\text{pz-of-alt1} :: ('a, 'k) \text{ alt} \Rightarrow ('a \text{ option}, 'k1, 'k) \text{ pz}$ **where**

[simp]: $\text{pz-of-alt1} = \text{unfold-pz} (\text{pz-of-alt1-emb} \circ \text{un-Alt})$

definition $\text{pz-of-alt2} :: ('a, 'k) \text{ alt} \Rightarrow ('a \text{ option}, 'k, 'k2) \text{ pz}$ **where**

[simp]: $\text{pz-of-alt2} = \text{unfold-pz} (\text{pz-of-alt2-emb} \circ \text{un-Alt})$

5 Automation Setup

lemma $\text{mc-rel-eq}[\text{unfolded vimage2p-def}]$:

$\text{BNF-Def.vimage2p un-MC un-MC (rel-pmf (=))} = (=)$

$\langle \text{proof} \rangle$

lemma $\text{dlts-rel-eq}[\text{unfolded vimage2p-def}]$:

$\text{BNF-Def.vimage2p un-DLTS un-DLTS (rel-fun (=) (rel-option (=)))} = (=)$

$\langle \text{proof} \rangle$

lemma $\text{react-rel-eq}[\text{unfolded vimage2p-def}]$:

$\text{BNF-Def.vimage2p un-React un-React (rel-fun (=) (rel-option (rel-pmf (=))))} = (=)$

$\langle \text{proof} \rangle$

lemma $\text{all-neq-Inl-ex-eq-Inr}[\text{dest}]$: $(\forall l. x \neq \text{Inl } l) \Longrightarrow (\exists r. x = \text{Inr } r)$ $\langle \text{proof} \rangle$

lemma $\text{all-neq-Inr-ex-eq-Inl}[\text{dest}]$: $(\forall r. x \neq \text{Inr } r) \Longrightarrow (\exists l. x = \text{Inl } l)$ $\langle \text{proof} \rangle$

lemma $\text{all2-neq-Inl-ex-eq-Inr}[\text{dest}]$: $(\forall a b. x \neq \text{Inl } (a, b)) \Longrightarrow (\exists r. x = \text{Inr } r)$ $\langle \text{proof} \rangle$

lemma $\text{all2-neq-Inr-ex-eq-Inl}[\text{dest}]$: $(\forall a b. x \neq \text{Inr } (a, b)) \Longrightarrow (\exists l. x = \text{Inl } l)$ $\langle \text{proof} \rangle$

lemma $\text{rel-prod-simp-asym}[\text{simp}]$:

$\bigwedge x y. \text{rel-prod } R S (x, y) = (\lambda z. \text{case } z \text{ of } (x', y') \Rightarrow R x x' \wedge S y y')$

$\bigwedge x y z. \text{rel-prod } R \ S \ x \ (y, z) = (\text{case } x \text{ of } (y', z') \Rightarrow R \ y' \ y \wedge S \ z' \ z)$
 <proof>

lemma *map-prod-eq-Pair-iff*[simp]:
 $\text{map-prod } f \ g \ x = (y, z) \longleftrightarrow (f \ (\text{fst } x) = y \wedge g \ (\text{snd } x) = z)$
 <proof>

lemmas [abs-def, simp] =
 sum.rel-map prod.rel-map option.rel-map pmf.rel-map bset.rel-map fun.rel-map
 nebset.rel-map

lemmas [simp] =
 lts.rel-eq lmc.rel-eq lmdp.rel-eq gen.rel-eq str.rel-eq alt.rel-eq sseq.rel-eq seg.rel-eq
 bun.rel-eq pz.rel-eq mg.rel-eq
 rel-pmf-return-pmf1 rel-pmf-return-pmf2 set-pmf-not-empty rel-pmf-rel-prod
 bset.set-map nebset.set-map

lemmas [simp del] =
 split-paired-Ex

lemma *bisimilar-eqI*:
assumes $\bigwedge R. \llbracket R \ x \ y; \bigwedge x y. R \ x \ y \Longrightarrow Q \ R \ (s1 \ x) \ (s2 \ y) \rrbracket \Longrightarrow P \ x \ y$
and $P \ x \ y \Longrightarrow \forall x y. P \ x \ y \longrightarrow Q \ P \ (s1 \ x) \ (s2 \ y)$
shows *bisimilar* $Q \ s1 \ s2 \ x \ y = P \ x \ y$
 <proof>

bundle *probabilistic-hierarchy* =
 rel-fun-def[simp]
 sum.splits[split]
 prod.splits[split]
 option.splits[split]

predicate2-eqD[THEN iffD2, OF mc-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF dlts-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF lts-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF react-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF lmc-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF lmdp-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF gen-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF str-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF alt-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF sseq-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF seg-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF bun-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF pz-rel-eq, dest]
predicate2-eqD[THEN iffD2, OF mg-rel-eq, dest]

iffD1[OF lts-rel-sel, dest!]
iffD1[OF lmc-rel-sel, dest!]

iffD1 [OF lmdp.rel-sel, dest!]
 iffD1 [OF gen.rel-sel, dest!]
 iffD1 [OF str.rel-sel, dest!]
 iffD1 [OF alt.rel-sel, dest!]
 iffD1 [OF sseg.rel-sel, dest!]
 iffD1 [OF seg.rel-sel, dest!]
 iffD1 [OF bun.rel-sel, dest!]
 iffD1 [OF pz.rel-sel, dest!]
 iffD1 [OF mg.rel-sel, dest!]

pmf.rel-refl [intro]
 bset.rel-refl [intro]
 nebset.rel-refl [intro]
 prod.rel-refl [intro]
 sum.rel-refl [intro]
 option.rel-refl [intro]

$\text{pmf.rel-mono-strong}$ [intro]
 $\text{bset.rel-mono-strong}$ [intro]
 $\text{nebset.rel-mono-strong}$ [intro]
 $\text{prod.rel-mono-strong}$ [intro]
 $\text{sum.rel-mono-strong}$ [intro]
 $\text{option.rel-mono-strong}$ [intro]

6 Proofs

context
includes *probabilistic-hierarchy*
begin

method *bisimilar-alt* =
rule *bisimilar-eqI*,
match conclusion in $u1\ s1\ x = u2\ s2\ y$ **for** $u1\ u2\ s1\ s2\ x\ y \Rightarrow$
<coinduction arbitrary: x y, fastforce>,
fastforce

lemma *bisimilar-alt*:

$\bigwedge s1\ s2. \text{bisimilar-mc } s1\ s2\ x\ y = (\text{unfold-mc } s1\ x = \text{unfold-mc } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-dlts } s1\ s2\ x\ y = (\text{unfold-dlts } s1\ x = \text{unfold-dlts } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-lts } s1\ s2\ x\ y = (\text{unfold-lts } s1\ x = \text{unfold-lts } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-react } s1\ s2\ x\ y = (\text{unfold-react } s1\ x = \text{unfold-react } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-lmc } s1\ s2\ x\ y = (\text{unfold-lmc } s1\ x = \text{unfold-lmc } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-lmdp } s1\ s2\ x\ y = (\text{unfold-lmdp } s1\ x = \text{unfold-lmdp } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-gen } s1\ s2\ x\ y = (\text{unfold-gen } s1\ x = \text{unfold-gen } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-str } s1\ s2\ x\ y = (\text{unfold-str } s1\ x = \text{unfold-str } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y = (\text{unfold-alt } s1\ x = \text{unfold-alt } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-sseg } s1\ s2\ x\ y = (\text{unfold-sseg } s1\ x = \text{unfold-sseg } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-seg } s1\ s2\ x\ y = (\text{unfold-seg } s1\ x = \text{unfold-seg } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-bun } s1\ s2\ x\ y = (\text{unfold-bun } s1\ x = \text{unfold-bun } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-pz } s1\ s2\ x\ y = (\text{unfold-pz } s1\ x = \text{unfold-pz } s2\ y)$
 $\bigwedge s1\ s2. \text{bisimilar-mg } s1\ s2\ x\ y = (\text{unfold-mg } s1\ x = \text{unfold-mg } s2\ y)$
 ⟨proof⟩

method *commute-prover* =
 intro ext,
 match **conclusion** in $u1\ s1\ x = (\text{emb } o\ u2\ s2)\ x$ **for** $\text{emb } u1\ u2\ s1\ s2\ x \Rightarrow$
 ⟨coinduction arbitrary: x , fastforce⟩

lemma *emb-commute*:

$\bigwedge s. \text{unfold-lts } (\text{lts-of-dlts-emb } o\ s) = \text{lts-of-dlts } o\ \text{unfold-dlts } s$
 $\bigwedge s. \text{unfold-gen } (\text{gen-of-lmc-emb } o\ s) = \text{gen-of-lmc } o\ \text{unfold-lmc } s$
 $\bigwedge s. \text{unfold-lmdp } (\text{lmdp-of-lmc-emb } o\ s) = \text{lmdp-of-lmc } o\ \text{unfold-lmc } s$
 $\bigwedge s. \text{unfold-react } (\text{react-of-dlts-emb } o\ s) = \text{react-of-dlts } o\ \text{unfold-dlts } s$
 $\bigwedge s. \text{unfold-sseg } (\text{sseg-of-lmdp-emb } o\ s) = \text{sseg-of-lmdp } o\ \text{unfold-lmdp } s$
 $\bigwedge s. \text{unfold-sseg } (\text{sseg-of-lts-emb } o\ s) = \text{sseg-of-lts } o\ \text{unfold-lts } s$
 $\bigwedge s. \text{unfold-sseg } (\text{ssegopt-of-alt-emb } o\ s) = \text{ssegopt-of-alt } o\ \text{unfold-alt } s$
 $\bigwedge s. \text{unfold-sseg } (\text{sseg-of-react-emb } o\ s) = \text{sseg-of-react } o\ \text{unfold-react } s$
 $\bigwedge s. \text{unfold-seg } (\text{seg-of-sseg-emb } o\ s) = \text{seg-of-sseg } o\ \text{unfold-sseg } s$
 $\bigwedge s. \text{unfold-seg } (\text{seg-of-gen-emb } o\ s) = \text{seg-of-gen } o\ \text{unfold-gen } s$
 $\bigwedge s. \text{unfold-bun } (\text{bun-of-lts-emb } o\ s) = \text{bun-of-lts } o\ \text{unfold-lts } s$
 $\bigwedge s. \text{unfold-bun } (\text{bunopt-of-alt-emb } o\ s) = \text{bunopt-of-alt } o\ \text{unfold-alt } s$
 $\bigwedge s. \text{unfold-bun } (\text{bun-of-gen-emb } o\ s) = \text{bun-of-gen } o\ \text{unfold-gen } s$
 $\bigwedge s. \text{unfold-pz } (\text{pz-of-seg-emb } o\ s) = \text{pz-of-seg } o\ \text{unfold-seg } s$
 $\bigwedge s. \text{unfold-pz } (\text{pz-of-bun-emb } o\ s) = \text{pz-of-bun } o\ \text{unfold-bun } s$
 $\bigwedge s. \text{unfold-str } (\text{str-of-mc-emb } o\ s) = \text{str-of-mc } o\ \text{unfold-mc } s$
 $\bigwedge s. \text{unfold-alt } (\text{alt-of-str-emb } o\ s) = \text{alt-of-str } o\ \text{unfold-str } s$
 $\bigwedge s. \text{unfold-sseg } (\text{ssegopt-of-sseg-emb } o\ s) = \text{ssegopt-of-sseg } o\ \text{unfold-sseg } s$
 $\bigwedge s. \text{unfold-seg } (\text{segopt-of-seg-emb } o\ s) = \text{segopt-of-seg } o\ \text{unfold-seg } s$
 $\bigwedge s. \text{unfold-bun } (\text{bunopt-of-bun-emb } o\ s) = \text{bunopt-of-bun } o\ \text{unfold-bun } s$
 $\bigwedge s. \text{unfold-pz } (\text{pzopt-of-pz-emb } o\ s) = \text{pzopt-of-pz } o\ \text{unfold-pz } s$
 $\bigwedge s. \text{unfold-pz } (\text{pzopt-of-mg-emb } o\ s) = \text{pzopt-of-mg } o\ \text{unfold-mg } s$
 $\bigwedge s. \text{unfold-mg } (\text{mg-of-pzopt-emb } o\ s) = \text{mg-of-pzopt } o\ \text{unfold-pz } s$

 $\bigwedge s. \text{unfold-mg } (\text{mg-of-pz-emb } o\ s) = \text{mg-of-pz } o\ \text{unfold-pz } s$
 $\bigwedge s. \text{unfold-mg } (\text{mg-of-alt1-emb } o\ s) = \text{mg-of-alt1 } o\ \text{unfold-alt } s$
 $\bigwedge s. \text{unfold-mg } (\text{mg-of-alt2-emb } o\ s) = \text{mg-of-alt2 } o\ \text{unfold-alt } s$
 $\bigwedge s. \text{unfold-pz } (\text{pz-of-alt1-emb } o\ s) = \text{pz-of-alt1 } o\ \text{unfold-alt } s$
 $\bigwedge s. \text{unfold-pz } (\text{pz-of-alt2-emb } o\ s) = \text{pz-of-alt2 } o\ \text{unfold-alt } s$
 ⟨proof⟩

method *inj-prover* =
 intro injI,
 match **conclusion** in $x = y$ **for** $x\ y \Rightarrow$ ⟨coinduction arbitrary: $x\ y$, fastforce⟩

lemma *inj*:

inj lts-of-dlts
inj react-of-dlts
inj gen-of-lmc

inj lmdp-of-lmc
inj sseg-of-lmdp
inj sseg-of-react
inj sseg-of-lts
inj ssegopt-of-alt
inj seg-of-gen
inj seg-of-sseg
inj bun-of-lts
inj bunopt-of-alt
inj bun-of-gen
inj pz-of-seg
inj pz-of-bun
inj str-of-mc
inj alt-of-str
inj ssegopt-of-sseg
inj segopt-of-seg
inj bunopt-of-bun
inj pzopt-of-pz
inj pzopt-of-mg
inj mg-of-pzopt

inj mg-of-pz
inj mg-of-alt1
inj mg-of-alt2
inj pz-of-alt1
inj pz-of-alt2
<proof>

end

lemma hierarchy:

$\bigwedge s1\ s2. \text{bisimilar-dlts } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-lts } (\text{lts-of-dlts-emb } o\ s1) (\text{lts-of-dlts-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lmc } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-gen } (\text{gen-of-lmc-emb } o\ s1) (\text{gen-of-lmc-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lmc } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-lmdp } (\text{lmdp-of-lmc-emb } o\ s1) (\text{lmdp-of-lmc-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-dlts } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-react } (\text{react-of-dlts-emb } o\ s1) (\text{react-of-dlts-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lmdp } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{sseg-of-lmdp-emb } o\ s1) (\text{sseg-of-lmdp-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lts } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{sseg-of-lts-emb } o\ s1) (\text{sseg-of-lts-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{ssegopt-of-alt-emb } o\ s1) (\text{ssegopt-of-alt-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-react } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{sseg-of-react-emb } o\ s1) (\text{sseg-of-react-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-sseg } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-seg } (\text{seg-of-sseg-emb } o\ s1) (\text{seg-of-sseg-emb } o\ s2)\ x\ y$

$\bigwedge (s1 :: - \Rightarrow ('a \text{ option} \times - \text{ pmf}) \text{ set}[-]) s2.$
 $\text{bisimilar-sseg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (\text{seg-of-sseg-emb } o \ s1) (\text{seg-of-sseg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-gen } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (\text{seg-of-gen-emb } o \ s1) (\text{seg-of-gen-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-lts } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (\text{bun-of-lts-emb } o \ s1) (\text{bun-of-lts-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-alt } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (\text{bunopt-of-alt-emb } o \ s1) (\text{bunopt-of-alt-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-gen } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (\text{bun-of-gen-emb } o \ s1) (\text{bun-of-gen-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-seg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (\text{pz-of-seg-emb } o \ s1) (\text{pz-of-seg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-bun } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (\text{pz-of-bun-emb } o \ s1) (\text{pz-of-bun-emb } o \ s2) \ x \ y$
 $\bigwedge (s1 :: - \Rightarrow ('a \text{ option} \times -) \text{ pmf } \text{ set}[-]) s2.$
 $\text{bisimilar-seg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (\text{pz-of-seg-emb } o \ s1) (\text{pz-of-seg-emb } o \ s2) \ x \ y$
 $\bigwedge (s1 :: - \Rightarrow (('a \text{ option} \times -) \text{ set}[-]) \text{ pmf}) s2.$
 $\text{bisimilar-bun } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (\text{pz-of-bun-emb } o \ s1) (\text{pz-of-bun-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-mc } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-str } (\text{str-of-mc-emb } o \ s1) (\text{str-of-mc-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-sseg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-sseg } (\text{ssegopt-of-sseg-emb } o \ s1) (\text{ssegopt-of-sseg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-seg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (\text{segopt-of-seg-emb } o \ s1) (\text{segopt-of-seg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-bun } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (\text{bunopt-of-bun-emb } o \ s1) (\text{bunopt-of-bun-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-pz } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (\text{pzopt-of-pz-emb } o \ s1) (\text{pzopt-of-pz-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-str } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-alt } (\text{alt-of-str-emb } o \ s1) (\text{alt-of-str-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-mg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (\text{pzopt-of-mg-emb } o \ s1) (\text{pzopt-of-mg-emb } o \ s2) \ x \ y$
 $\langle \text{proof} \rangle$

An edge that would make the graph cyclic

lemma

$\bigwedge s1 \ s2. \text{bisimilar-pz } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-mg } (\text{mg-of-pz-emb } o \ s1) (\text{mg-of-pz-emb } o \ s2) \ x \ y$
 $\langle \text{proof} \rangle$

Some redundant (historic) transitive edges

lemma

$\bigwedge s1 \ s2. \text{bisimilar-pz } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-mg } (\text{mg-of-pzopt-emb } o \ s1) (\text{mg-of-pzopt-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-alt } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-mg } (\text{mg-of-alt1-emb } o \ s1) (\text{mg-of-alt1-emb } o \ s2) \ x \ y$

$\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \iff \text{bisimilar-mg } (\text{mg-of-alt2-emb } o\ s1)\ (\text{mg-of-alt2-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \iff \text{bisimilar-pz } (\text{pz-of-alt1-emb } o\ s1)\ (\text{pz-of-alt1-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \iff \text{bisimilar-pz } (\text{pz-of-alt2-emb } o\ s1)\ (\text{pz-of-alt2-emb } o\ s2)\ x\ y$
 <proof>

7 Some special proofs

Two views on LTS

lemma $\exists f::((\text{'a} \times \text{'s})\ \text{set} \Rightarrow \text{'a} \Rightarrow \text{'s}\ \text{set}).\ \text{bij } f$
 <proof>

lemma $\exists f::((\text{'a} \times \text{'s})\ \text{set}[(\text{'a} \times \text{'s})\ \text{set}] \Rightarrow \text{'a} \Rightarrow \text{'s}\ \text{set}[\text{'s}\ \text{set}]).\ \text{bij } f$
 <proof>

mc is trivial

lemma *mc-unit*:

fixes $x\ y :: mc$

shows $x = y$

<proof>

lemma *bisimilar-mc* $s1\ s2\ x\ y$
 <proof>

8 Printing the Hierarchy Graph

<ML>

9 Vardi Systems

context notes $[[\text{bnf-internals}]]$

begin

datatype $(\text{'a}, \text{'b}, \text{'k})\ \text{var0} = \text{PMF } (\text{'a} \times \text{'b})\ \text{pmf} \mid \text{BPS } (\text{'a} \times \text{'b})\ \text{set}[\text{'k}]$

end

inductive *var-eq* $:: (\text{'a}, \text{'b}, \text{'k})\ \text{var0} \Rightarrow (\text{'a}, \text{'b}, \text{'k})\ \text{var0} \Rightarrow \text{bool}$ (**infixl** $\langle \sim \rangle$ 65)

where

var-eq-reflp[*intro*]: $x \sim x$

| [*intro*]: $\text{PMF } (\text{return-pmf } (a, x)) \sim \text{BPS } (\text{bsingleton } (a, x))$

| [*intro*]: $\text{BPS } (\text{bsingleton } (a, x)) \sim \text{PMF } (\text{return-pmf } (a, x))$

lemma *var-eq-symp*: $x \sim y \implies y \sim x$

<proof>

lemma *var-eq-transp*: $x \sim y \implies y \sim z \implies x \sim z$
 ⟨proof⟩

quotient-type $(\text{'a}, \text{'b}, \text{'k}) \text{ var} = (\text{'a}, \text{'b}, \text{'k}) \text{ var0} / (\sim)$
 ⟨proof⟩

lift-definition *map-var* :: $(\text{'a} \Rightarrow \text{'c}) \Rightarrow (\text{'b} \Rightarrow \text{'d}) \Rightarrow (\text{'a}, \text{'b}, \text{'k}) \text{ var} \Rightarrow (\text{'c}, \text{'d}, \text{'k}) \text{ var}$
is *map-var0*
 ⟨proof⟩

lift-definition *set1-var* :: $(\text{'a}, \text{'b}, \text{'k}) \text{ var} \Rightarrow \text{'a} \text{ set}$
is *set1-var0*
 ⟨proof⟩

lift-definition *set2-var* :: $(\text{'a}, \text{'b}, \text{'k}) \text{ var} \Rightarrow \text{'b} \text{ set}$
is *set2-var0*
 ⟨proof⟩

inductive *rel-var* :: $(\text{'a} \Rightarrow \text{'c} \Rightarrow \text{bool}) \Rightarrow (\text{'b} \Rightarrow \text{'d} \Rightarrow \text{bool}) \Rightarrow (\text{'a}, \text{'b}, \text{'k}) \text{ var} \Rightarrow (\text{'c}, \text{'d}, \text{'k}) \text{ var} \Rightarrow \text{bool}$ **for** *R S* **where**
set1-var $x \subseteq \{(x, y). R \ x \ y\} \implies \text{set2-var } x \subseteq \{(x, y). S \ x \ y\} \implies$
rel-var *R S* $(\text{map-var } \text{fst } \text{fst } x) (\text{map-var } \text{snd } \text{snd } x)$

abbreviation $(\text{input}) \text{ var0-of-gen-emb} \equiv \text{case-option } (\text{BPS } \text{bempty}) \text{ PMF}$
abbreviation $(\text{input}) \text{ var0-of-lts-emb} \equiv \text{BPS}$

lift-definition *var-of-gen-emb* :: $(\text{'a} \times \text{'b}) \text{ pmf option} \Rightarrow (\text{'a}, \text{'b}, \text{'k}) \text{ var}$ **is** *var0-of-gen-emb*
 ⟨proof⟩

lift-definition *var-of-lts-emb* :: $(\text{'a} \times \text{'b}) \text{ set['k]} \Rightarrow (\text{'a}, \text{'b}, \text{'k}) \text{ var}$ **is** *var0-of-lts-emb*
 ⟨proof⟩

abbreviation *bisimilar-var* $\equiv \text{bisimilar } (\lambda R. \text{rel-var } (=) R)$

lemma *map-var0-eq-BPS-iff[simp]*:
 $\text{map-var0 } f \ g \ z = \text{BPS } X \iff (\exists Y. z = \text{BPS } Y \wedge \text{map-bset } (\text{map-prod } f \ g) \ Y = X)$
 ⟨proof⟩

lemma *map-var0-eq-PMF-iff[simp]*:
 $\text{map-var0 } f \ g \ z = \text{PMF } p \iff (\exists q. z = \text{PMF } q \wedge \text{map-pmf } (\text{map-prod } f \ g) \ q = p)$
 ⟨proof⟩

lemma *bisimilar-lts* $s1 \ s2 \ x \ y \iff \text{bisimilar-var } (\text{var-of-lts-emb } o \ s1) (\text{var-of-lts-emb } o \ s2) \ x \ y$
(is $- \iff \text{bisimilar-var } (?emb1 \ o \ -) (?emb2 \ o \ -) \ -)$
 ⟨proof⟩

lemma *bisimilar-gen s1 s2 x y* \longleftrightarrow *bisimilar-var (var-of-gen-emb o s1) (var-of-gen-emb o s2) x y*
(**is** - \longleftrightarrow *bisimilar-var (?emb1 o -) (?emb2 o -) -*)
(*proof*)