

Probabilistic Hierarchy

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1 Bisimilarity

definition *bisimilar* where

$$\textit{bisimilar } Q \textit{ } s1 \textit{ } s2 \textit{ } x \textit{ } y \equiv (\exists R. R \textit{ } x \textit{ } y \wedge (\forall x \textit{ } y. R \textit{ } x \textit{ } y \longrightarrow Q \textit{ } R \textit{ } (s1 \textit{ } x) \textit{ } (s2 \textit{ } y)))$$

abbreviation *bisimilar-mc* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-pmf } R)$

abbreviation *bisimilar-dlts* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-fun } (=) \textit{ (rel-option } R))$

abbreviation *bisimilar-lts* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-bset } (\textit{rel-prod } (=) \textit{ } R))$

abbreviation *bisimilar-react* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-fun } (=) \textit{ (rel-option } (\textit{rel-pmf } R)))$

abbreviation *bisimilar-lmc* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-prod } (=) \textit{ (rel-pmf } R))$

abbreviation *bisimilar-lmdp* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-prod } (=) \textit{ (rel-nebset } (\textit{rel-pmf } R)))$

abbreviation *bisimilar-gen* $\equiv \textit{bisimilar } (\lambda R. \textit{rel-option } (\textit{rel-pmf } (\textit{rel-prod } (=) \textit{ } R)))$

abbreviation $\text{bisimilar-str} \equiv \text{bisimilar } (\lambda R. \text{rel-sum } (\text{rel-pmf } R) (\text{rel-option } (\text{rel-prod } (=) R)))$
abbreviation $\text{bisimilar-alt} \equiv \text{bisimilar } (\lambda R. \text{rel-sum } (\text{rel-pmf } R) (\text{rel-bset } (\text{rel-prod } (=) R)))$
abbreviation $\text{bisimilar-sseg} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-prod } (=) (\text{rel-pmf } R)))$
abbreviation $\text{bisimilar-seg} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-pmf } (\text{rel-prod } (=) R)))$
abbreviation $\text{bisimilar-bun} \equiv \text{bisimilar } (\lambda R. \text{rel-pmf } (\text{rel-bset } (\text{rel-prod } (=) R)))$
abbreviation $\text{bisimilar-pz} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-pmf } (\text{rel-bset } (\text{rel-prod } (=) R))))$
abbreviation $\text{bisimilar-mg} \equiv \text{bisimilar } (\lambda R. \text{rel-bset } (\text{rel-pmf } (\text{rel-bset } (\text{rel-sum } (\text{rel-prod } (=) R) R))))$

2 Systems

codatatype $\text{mc} = \text{MC } \text{mc } \text{pmf}$
codatatype $'a \text{ dlts} = \text{DLTS } 'a \Rightarrow 'a \text{ dlts } \text{option}$
codatatype $('a, 'k) \text{ lts} = \text{LTS } ('a \times ('a, 'k) \text{ lts}) \text{ set}['k]$
codatatype $'a \text{ react} = \text{React } 'a \Rightarrow 'a \text{ react } \text{pmf } \text{option}$
codatatype $'a \text{ lmc} = \text{LMC } 'a \times 'a \text{ lmc } \text{pmf}$
codatatype $('a, 'k) \text{ lmdp} = \text{LMDP } 'a \times ('a, 'k) \text{ lmdp } \text{pmf } \text{set!}['k]$
codatatype $'a \text{ gen} = \text{Gen } ('a \times 'a \text{ gen}) \text{pmf } \text{option}$
codatatype $'a \text{ str} = \text{Str } 'a \text{ str } \text{pmf} + ('a \times 'a \text{ str}) \text{option}$
codatatype $('a, 'k) \text{ alt} = \text{Alt } ('a, 'k) \text{ alt } \text{pmf} + ('a \times ('a, 'k) \text{ alt}) \text{set}['k]$
codatatype $('a, 'k) \text{ sseg} = \text{SSeg } ('a \times ('a, 'k) \text{ sseg } \text{pmf}) \text{set}['k]$
codatatype $('a, 'k) \text{ seg} = \text{Seg } ('a \times ('a, 'k) \text{ seg}) \text{pmf } \text{set}['k]$
codatatype $('a, 'k) \text{ bun} = \text{Bun } (('a \times ('a, 'k) \text{ bun}) \text{set}['k]) \text{pmf}$
codatatype $('a, 'k1, 'k2) \text{ pz} = \text{PZ } (('a \times ('a, 'k1, 'k2) \text{ pz}) \text{set}['k1]) \text{pmf } \text{set}['k2]$
codatatype $('a, 'k1, 'k2) \text{ mg} = \text{MG } (('a \times ('a, 'k1, 'k2) \text{ mg} + ('a, 'k1, 'k2) \text{ mg}) \text{set}['k1]) \text{pmf } \text{set}['k2]$

3 Unfolds

primcorec $\text{unfold-mc} :: ('a \Rightarrow 'a \text{ pmf}) \Rightarrow 'a \Rightarrow \text{mc } \mathbf{where}$
 $\text{unfold-mc } s \ x = \text{MC } (\text{map-pmf } (\text{unfold-mc } s) (s \ x))$

primcorec $\text{unfold-dlts} :: ('a \Rightarrow 'b \Rightarrow 'a \text{ option}) \Rightarrow 'a \Rightarrow 'b \text{ dlts } \mathbf{where}$
 $\text{unfold-dlts } s \ x = \text{DLTS } (\text{map-option } (\text{unfold-dlts } s) o \ s \ x)$

primcorec $\text{unfold-lts} :: ('a \Rightarrow ('b \times 'a) \text{ set}['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ lts } \mathbf{where}$
 $\text{unfold-lts } s \ x = \text{LTS } (\text{map-bset } (\text{map-prod } \text{id } (\text{unfold-lts } s)) (s \ x))$

primcorec $\text{unfold-react} :: ('a \Rightarrow 'b \Rightarrow 'a \text{ pmf } \text{option}) \Rightarrow 'a \Rightarrow 'b \text{ react } \mathbf{where}$
 $\text{unfold-react } s \ x = \text{React } (\text{map-option } (\text{map-pmf } (\text{unfold-react } s)) o \ s \ x)$

primcorec $\text{unfold-lmc} :: ('a \Rightarrow 'b \times 'a \text{ pmf}) \Rightarrow 'a \Rightarrow 'b \text{ lmc } \mathbf{where}$
 $\text{unfold-lmc } s \ x = \text{LMC } (\text{map-prod } \text{id } (\text{map-pmf } (\text{unfold-lmc } s)) (s \ x))$

primcorec $\text{unfold-lmdp} :: ('a \Rightarrow 'b \times 'a \text{ pmf } \text{set!}['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ lmdp } \mathbf{where}$

unfold-lmdp $s x = LMDP (map-prod id (map-nebset (map-pmf (unfold-lmdp s))) (s x))$

primcorec *unfold-gen* $:: ('a \Rightarrow (('b \times 'a) pmf) option) \Rightarrow 'a \Rightarrow 'b \text{ gen where}$
unfold-gen $s x = Gen (map-option (map-pmf (map-prod id (unfold-gen s))) (s x))$

primcorec *unfold-str* $:: ('a \Rightarrow 'a pmf + ('b \times 'a) option) \Rightarrow 'a \Rightarrow 'b \text{ str where}$
unfold-str $s x = Str (map-sum (map-pmf (unfold-str s)) (map-option (map-prod id (unfold-str s)))) (s x)$

primcorec *unfold-alt* $:: ('a \Rightarrow 'a pmf + ('b \times 'a) set['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ alt where}$
unfold-alt $s x = Alt (map-sum (map-pmf (unfold-alt s)) (map-bset (map-prod id (unfold-alt s)))) (s x)$

primcorec *unfold-sseg* $:: ('a \Rightarrow ('b \times 'a pmf) set['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ sseg where}$
unfold-sseg $s x = SSeg (map-bset (map-prod id (map-pmf (unfold-sseg s)))) (s x)$

primcorec *unfold-seg* $:: ('a \Rightarrow (('b \times 'a) pmf) set['k]) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ seg where}$
unfold-seg $s x = Seg (map-bset (map-pmf (map-prod id (unfold-seg s)))) (s x)$

primcorec *unfold-bun* $:: ('a \Rightarrow (('b \times 'a) set['k]) pmf) \Rightarrow 'a \Rightarrow ('b, 'k) \text{ bun where}$
unfold-bun $s x = Bun (map-pmf (map-bset (map-prod id (unfold-bun s)))) (s x)$

primcorec *unfold-pz* $:: ('a \Rightarrow (('b \times 'a) set['k1]) pmf set['k2]) \Rightarrow 'a \Rightarrow ('b, 'k1, 'k2) \text{ pz where}$
unfold-pz $s x = PZ (map-bset (map-pmf (map-bset (map-prod id (unfold-pz s)))) (s x))$

primcorec *unfold-mg* $:: ('a \Rightarrow (('b \times 'a + 'a) set['k1]) pmf set['k2]) \Rightarrow 'a \Rightarrow ('b, 'k1, 'k2) \text{ mg where}$
unfold-mg $s x = MG (map-bset (map-pmf (map-bset (map-sum (map-prod id (unfold-mg s)) (unfold-mg s)))) (s x))$

4 Embeddings

abbreviation (*input*) *react-of-dlts-emb* $dlts \equiv map-option return-pmf o dlts$

abbreviation (*input*) *lts-of-dlts-emb* $\equiv bgraph$

abbreviation (*input*) *sseg-of-react-emb* $\equiv bgraph$

abbreviation (*input*) *gen-of-lmc-emb* $\equiv Some o case-prod (map-pmf o Pair)$

abbreviation (*input*) *lmdp-of-lmc-emb* $\equiv map-prod id nebsingleton$

abbreviation (*input*) *sseg-of-lmdp-emb* $\equiv (\lambda(a, X). map-bset (Pair a) (bset-of-nebset X))$

abbreviation (*input*) *sseg-of-lts-emb* $\equiv map-bset (map-prod id return-pmf)$

abbreviation (*input*) *ssegopt-of-alt-emb* $\equiv case-sum$

$(map-bset (Pair None) o bsingleton)$

$(map-bset (map-prod Some return-pmf))$

abbreviation (*input*) *bunopt-of-alt-emb* $\equiv case-sum$

$(\text{map-pmf } (\text{bsingleton } o \text{ Pair None}))$
 $(\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id)) o \text{ return-pmf})$
abbreviation $(\text{input}) \text{ segopt-of-seg-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-prod } \text{Some } id))$
abbreviation $(\text{input}) \text{ ssegopt-of-sseg-emb} \equiv \text{map-bset } (\text{map-prod } \text{Some } id)$
abbreviation $(\text{input}) \text{ bunopt-of-bun-emb} \equiv \text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id))$
abbreviation $(\text{input}) \text{ pzopt-of-pz-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id)))$
abbreviation $(\text{input}) \text{ seg-of-sseg-emb} \equiv \text{map-bset } (\text{case-prod } (\text{map-pmf } o \text{ Pair}))$
abbreviation $(\text{input}) \text{ pz-of-seg-emb} \equiv \text{map-bset } (\text{map-pmf } \text{bsingleton})$
abbreviation $(\text{input}) \text{ pz-of-bun-emb} \equiv \text{bsingleton}$
abbreviation $(\text{input}) \text{ seg-of-gen-emb} \equiv \text{bset-of-option}$
abbreviation $(\text{input}) \text{ bun-of-lts-emb} \equiv \text{return-pmf}$
abbreviation $(\text{input}) \text{ bun-of-gen-emb} \equiv \text{case-option } (\text{return-pmf } \text{bempty}) (\text{map-pmf } \text{bsingleton})$
abbreviation $(\text{input}) \text{ str-of-mc-emb} \equiv \text{Inl}$
abbreviation $(\text{input}) \text{ alt-of-str-emb} \equiv \text{map-sum } id \text{ bset-of-option}$
abbreviation $(\text{input}) \text{ pzopt-of-mg-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{case-sum } (\text{map-prod } \text{Some } id) (\text{Pair } \text{None}))))$
abbreviation $(\text{input}) \text{ mg-of-pzopt-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } (\lambda(a, s). \text{case-option } (\text{Inr } s) (\lambda a. (\text{Inl } (a, s))) a))))$

Obsolete edges (susumed by transitive ones)

abbreviation $(\text{input}) \text{ mg-of-pz-emb} \equiv \text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inl}))$
abbreviation $(\text{input}) \text{ mg-of-alt1-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inr } o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inl } o \text{ bsingleton}) o \text{ return-pmf}))$
abbreviation $(\text{input}) \text{ mg-of-alt2-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inr } o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } \text{Inl } o \text{ return-pmf}) o \text{ bsingleton}))$
abbreviation $(\text{input}) \text{ pz-of-alt1-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{Pair } \text{None}) o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id) o \text{ bsingleton}) o \text{ return-pmf}))$
abbreviation $(\text{input}) \text{ pz-of-alt2-emb} \equiv \text{case-sum } (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{Pair } \text{None}) o \text{ bsingleton})) o \text{ bsingleton}) (\text{map-bset } (\text{map-pmf } (\text{map-bset } (\text{map-prod } \text{Some } id)) o \text{ return-pmf}) o \text{ bsingleton})$

definition $\text{react-of-dlts} :: 'a \text{ dlts} \Rightarrow 'a \text{ react}$ **where**
 $[\text{simp}]: \text{react-of-dlts} = \text{unfold-react } (\text{react-of-dlts-emb } o \text{ un-DLTS})$

definition $\text{lts-of-dlts} :: 'a \text{ dlts} \Rightarrow ('a, 'a \text{ set}) \text{ lts}$ **where**
 $[\text{simp}]: \text{lts-of-dlts} = \text{unfold-lts } (\text{lts-of-dlts-emb } o \text{ un-DLTS})$

definition $\text{sseg-of-react} :: 'a \text{ react} \Rightarrow ('a, 'a \text{ set}) \text{ sseg}$ **where**
 $[\text{simp}]: \text{sseg-of-react} = \text{unfold-sseg } (\text{sseg-of-react-emb } o \text{ un-React})$

definition $\text{lmdp-of-lmc} :: 'a \text{ lmc} \Rightarrow ('a, 'k) \text{ lmdp}$ **where**

[simp]: $lmdp\text{-of-lmc} = \text{unfold-lmdp } (lmdp\text{-of-lmc-emb } o \text{ un-LMC})$

definition $gen\text{-of-lmc} :: 'a \text{ lmc} \Rightarrow 'a \text{ gen}$ **where**

[simp]: $gen\text{-of-lmc} = \text{unfold-gen } (gen\text{-of-lmc-emb } o \text{ un-LMC})$

definition $sseg\text{-of-lmdp} :: ('a, 'k) \text{ lmdp} \Rightarrow ('a, 'k) \text{ sseg}$ **where**

[simp]: $sseg\text{-of-lmdp} = \text{unfold-sseg } (sseg\text{-of-lmdp-emb } o \text{ un-LMDP})$

definition $sseg\text{-of-lts} :: ('a, 'k) \text{ lts} \Rightarrow ('a, 'k) \text{ sseg}$ **where**

[simp]: $sseg\text{-of-lts} = \text{unfold-sseg } (sseg\text{-of-lts-emb } o \text{ un-LTS})$

definition $ssegopt\text{-of-alt} :: ('a, 'k) \text{ alt} \Rightarrow ('a \text{ option}, 'k) \text{ sseg}$ **where**

[simp]: $ssegopt\text{-of-alt} = \text{unfold-sseg } (ssegopt\text{-of-alt-emb } o \text{ un-Alt})$

definition $bunopt\text{-of-alt} :: ('a, 'k) \text{ alt} \Rightarrow ('a \text{ option}, 'k) \text{ bun}$ **where**

[simp]: $bunopt\text{-of-alt} = \text{unfold-bun } (bunopt\text{-of-alt-emb } o \text{ un-Alt})$

definition $seg\text{-of-sseg} :: ('a, 'k) \text{ sseg} \Rightarrow ('a, 'k) \text{ seg}$ **where**

[simp]: $seg\text{-of-sseg} = \text{unfold-seg } (seg\text{-of-sseg-emb } o \text{ un-SSeg})$

definition $seg\text{-of-gen} :: 'a \text{ gen} \Rightarrow ('a, 'k) \text{ seg}$ **where**

[simp]: $seg\text{-of-gen} = \text{unfold-seg } (seg\text{-of-gen-emb } o \text{ un-Gen})$

definition $bun\text{-of-lts} :: ('a, 'k) \text{ lts} \Rightarrow ('a, 'k) \text{ bun}$ **where**

[simp]: $bun\text{-of-lts} = \text{unfold-bun } (bun\text{-of-lts-emb } o \text{ un-LTS})$

definition $bun\text{-of-gen} :: 'a \text{ gen} \Rightarrow ('a, 'k) \text{ bun}$ **where**

[simp]: $bun\text{-of-gen} = \text{unfold-bun } (bun\text{-of-gen-emb } o \text{ un-Gen})$

definition $pz\text{-of-seg} :: ('a, 'k) \text{ seg} \Rightarrow ('a, 'k1, 'k) \text{ pz}$ **where**

[simp]: $pz\text{-of-seg} = \text{unfold-pz } (pz\text{-of-seg-emb } o \text{ un-Seg})$

definition $pz\text{-of-bun} :: ('a, 'k) \text{ bun} \Rightarrow ('a, 'k, 'k1) \text{ pz}$ **where**

[simp]: $pz\text{-of-bun} = \text{unfold-pz } (pz\text{-of-bun-emb } o \text{ un-Bun})$

definition $mg\text{-of-pz} :: ('a, 'k1, 'k2) \text{ pz} \Rightarrow ('a, 'k1, 'k2) \text{ mg}$ **where**

[simp]: $mg\text{-of-pz} = \text{unfold-mg } (mg\text{-of-pz-emb } o \text{ un-PZ})$

definition $str\text{-of-mc} :: mc \Rightarrow 'a \text{ str}$ **where**

[simp]: $str\text{-of-mc} = \text{unfold-str } (str\text{-of-mc-emb } o \text{ un-MC})$

definition $alt\text{-of-str} :: 'a \text{ str} \Rightarrow ('a, 'k) \text{ alt}$ **where**

[simp]: $alt\text{-of-str} = \text{unfold-alt } (alt\text{-of-str-emb } o \text{ un-Str})$

definition $ssegopt\text{-of-sseg} :: ('a, 'k) \text{ sseg} \Rightarrow ('a \text{ option}, 'k) \text{ sseg}$ **where**

[simp]: $ssegopt\text{-of-sseg} = \text{unfold-sseg } (ssegopt\text{-of-sseg-emb } o \text{ un-SSeg})$

definition $segopt\text{-of-seg} :: ('a, 'k) \text{ seg} \Rightarrow ('a \text{ option}, 'k) \text{ seg}$ **where**

[simp]: $segopt\text{-of-seg} = \text{unfold-seg } (segopt\text{-of-seg-emb } o \text{ un-Seg})$

definition *bunopt-of-bun* :: ('a, 'k) bun \Rightarrow ('a option, 'k) bun **where**

[simp]: *bunopt-of-bun* = *unfold-bun* (*bunopt-of-bun-emb* o *un-Bun*)

definition *pzopt-of-pz* :: ('a, 'k1, 'k2) pz \Rightarrow ('a option, 'k1, 'k2) pz **where**

[simp]: *pzopt-of-pz* = *unfold-pz* (*pzopt-of-pz-emb* o *un-PZ*)

definition *pzopt-of-mg* :: ('a, 'k1, 'k2) mg \Rightarrow ('a option, 'k1, 'k2) pz **where**

[simp]: *pzopt-of-mg* = *unfold-pz* (*pzopt-of-mg-emb* o *un-MG*)

definition *mg-of-pzopt* :: ('a option, 'k1, 'k2) pz \Rightarrow ('a, 'k1, 'k2) mg **where**

[simp]: *mg-of-pzopt* = *unfold-mg* (*mg-of-pzopt-emb* o *un-PZ*)

definition *mg-of-alt1* :: ('a, 'k) alt \Rightarrow ('a, 'k1, 'k) mg **where**

[simp]: *mg-of-alt1* = *unfold-mg* (*mg-of-alt1-emb* o *un-Alt*)

definition *mg-of-alt2* :: ('a, 'k) alt \Rightarrow ('a, 'k, 'k1) mg **where**

[simp]: *mg-of-alt2* = *unfold-mg* (*mg-of-alt2-emb* o *un-Alt*)

definition *pz-of-alt1* :: ('a, 'k) alt \Rightarrow ('a option, 'k1, 'k) pz **where**

[simp]: *pz-of-alt1* = *unfold-pz* (*pz-of-alt1-emb* o *un-Alt*)

definition *pz-of-alt2* :: ('a, 'k) alt \Rightarrow ('a option, 'k, 'k2) pz **where**

[simp]: *pz-of-alt2* = *unfold-pz* (*pz-of-alt2-emb* o *un-Alt*)

5 Automation Setup

lemma *mc-rel-eq*[*unfolded vimage2p-def*]:

BNF-Def.vimage2p un-MC un-MC (rel-pmf (=)) = (=)

by (*auto simp add: vimage2p-def pmf.rel-eq option.rel-eq fun.rel-eq fun-eq-iff mc.expand*)

lemma *dlts-rel-eq*[*unfolded vimage2p-def*]:

BNF-Def.vimage2p un-DLTS un-DLTS (rel-fun (=) (rel-option (=))) = (=)

by (*auto simp add: vimage2p-def pmf.rel-eq option.rel-eq fun.rel-eq fun-eq-iff dlts.expand*)

lemma *react-rel-eq*[*unfolded vimage2p-def*]:

BNF-Def.vimage2p un-React un-React (rel-fun (=) (rel-option (rel-pmf (=)))) = (=)

by (*auto simp add: vimage2p-def pmf.rel-eq option.rel-eq fun.rel-eq fun-eq-iff react.expand*)

lemma *all-neq-Inl-ex-eq-Inr*[*dest*]: $(\forall l. x \neq \text{Inl } l) \implies (\exists r. x = \text{Inr } r)$ **by** (*cases x*) *auto*

lemma *all-neq-Inr-ex-eq-Inl*[*dest*]: $(\forall r. x \neq \text{Inr } r) \implies (\exists l. x = \text{Inl } l)$ **by** (*cases x*) *auto*

lemma *all2-neq-Inl-ex-eq-Inr*[*dest*]: $(\forall a b. x \neq \text{Inl } (a, b)) \implies (\exists r. x = \text{Inr } r)$ **by** (*cases x*) *auto*

lemma *all2-neq-Inr-ex-eq-Inl*[*dest*]: $(\forall a b. x \neq \text{Inr } (a, b)) \implies (\exists l. x = \text{Inl } l)$ **by**
(cases x) auto

lemma *rel-prod-simp-asym*[*simp*]:

$\bigwedge x y. \text{rel-prod } R S (x, y) = (\lambda z. \text{case } z \text{ of } (x', y') \Rightarrow R x x' \wedge S y y')$
 $\bigwedge x y z. \text{rel-prod } R S x (y, z) = (\text{case } x \text{ of } (y', z') \Rightarrow R y' y \wedge S z' z)$
by *auto*

lemma *map-prod-eq-Pair-iff*[*simp*]:

$\text{map-prod } f g x = (y, z) \iff (f (\text{fst } x) = y \wedge g (\text{snd } x) = z)$
by *(cases x) auto*

lemmas [*abs-def, simp*] =

sum.rel-map prod.rel-map option.rel-map pmf.rel-map bset.rel-map fun.rel-map
nebset.rel-map

lemmas [*simp*] =

lts.rel-eq lmc.rel-eq lmdp.rel-eq gen.rel-eq str.rel-eq alt.rel-eq sseq.rel-eq seg.rel-eq
bun.rel-eq pz.rel-eq mg.rel-eq
rel-pmf-return-pmf1 rel-pmf-return-pmf2 set-pmf-not-empty rel-pmf-rel-prod
bset.set-map nebset.set-map

lemmas [*simp del*] =

split-paired-Ex

lemma *bisimilar-eqI*:

assumes $\bigwedge R. \llbracket R x y; \bigwedge x y. R x y \implies Q R (s1 x) (s2 y) \rrbracket \implies P x y$
and $P x y \implies \forall x y. P x y \longrightarrow Q P (s1 x) (s2 y)$
shows *bisimilar* $Q s1 s2 x y = P x y$
using *assms unfolding bisimilar-def by auto*

bundle *probabilistic-hierarchy* =

rel-fun-def[*simp*]
sum.splits[*split*]
prod.splits[*split*]
option.splits[*split*]

predicate2-eqD[*THEN iffD2, OF mc-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF dlts-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF lts-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF react-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF lmc-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF lmdp-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF gen-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF str-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF alt-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF sseq-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF seg-rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF bun-rel-eq, dest*]

predicate2-eqD[*THEN iffD2, OF pz.rel-eq, dest*]
predicate2-eqD[*THEN iffD2, OF mg.rel-eq, dest*]

iffD1[*OF lts.rel-sel, dest!*]
iffD1[*OF lmc.rel-sel, dest!*]
iffD1[*OF lmdp.rel-sel, dest!*]
iffD1[*OF gen.rel-sel, dest!*]
iffD1[*OF str.rel-sel, dest!*]
iffD1[*OF alt.rel-sel, dest!*]
iffD1[*OF sseg.rel-sel, dest!*]
iffD1[*OF seg.rel-sel, dest!*]
iffD1[*OF bun.rel-sel, dest!*]
iffD1[*OF pz.rel-sel, dest!*]
iffD1[*OF mg.rel-sel, dest!*]

pmf.rel-refl[*intro*]
bset.rel-refl[*intro*]
nebset.rel-refl[*intro*]
prod.rel-refl[*intro*]
sum.rel-refl[*intro*]
option.rel-refl[*intro*]

pmf.rel-mono-strong[*intro*]
bset.rel-mono-strong[*intro*]
nebset.rel-mono-strong[*intro*]
prod.rel-mono-strong[*intro*]
sum.rel-mono-strong[*intro*]
option.rel-mono-strong[*intro*]

6 Proofs

context

includes *probabilistic-hierarchy*

begin

method *bisimilar-alt* =

rule bisimilar-eqI,

match conclusion in u1 s1 x = u2 s2 y for u1 u2 s1 s2 x y ⇒

⟨coinduction arbitrary: x y, fastforce⟩,

fastforce

lemma *bisimilar-alt*:

$\bigwedge s1\ s2. \text{bisimilar-mc } s1\ s2\ x\ y = (\text{unfold-mc } s1\ x = \text{unfold-mc } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-dlts } s1\ s2\ x\ y = (\text{unfold-dlts } s1\ x = \text{unfold-dlts } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-lts } s1\ s2\ x\ y = (\text{unfold-lts } s1\ x = \text{unfold-lts } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-react } s1\ s2\ x\ y = (\text{unfold-react } s1\ x = \text{unfold-react } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-lmc } s1\ s2\ x\ y = (\text{unfold-lmc } s1\ x = \text{unfold-lmc } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-lmdp } s1\ s2\ x\ y = (\text{unfold-lmdp } s1\ x = \text{unfold-lmdp } s2\ y)$

$\bigwedge s1\ s2. \text{bisimilar-gen } s1\ s2\ x\ y = (\text{unfold-gen } s1\ x = \text{unfold-gen } s2\ y)$

$\wedge s1\ s2. \text{bisimilar-str } s1\ s2\ x\ y = (\text{unfold-str } s1\ x = \text{unfold-str } s2\ y)$
 $\wedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y = (\text{unfold-alt } s1\ x = \text{unfold-alt } s2\ y)$
 $\wedge s1\ s2. \text{bisimilar-sseg } s1\ s2\ x\ y = (\text{unfold-sseg } s1\ x = \text{unfold-sseg } s2\ y)$
 $\wedge s1\ s2. \text{bisimilar-seg } s1\ s2\ x\ y = (\text{unfold-seg } s1\ x = \text{unfold-seg } s2\ y)$
 $\wedge s1\ s2. \text{bisimilar-bun } s1\ s2\ x\ y = (\text{unfold-bun } s1\ x = \text{unfold-bun } s2\ y)$
 $\wedge s1\ s2. \text{bisimilar-pz } s1\ s2\ x\ y = (\text{unfold-pz } s1\ x = \text{unfold-pz } s2\ y)$
 $\wedge s1\ s2. \text{bisimilar-mg } s1\ s2\ x\ y = (\text{unfold-mg } s1\ x = \text{unfold-mg } s2\ y)$
by *bisimilar-alt+*

method *commute-prover* =

intro ext,

match conclusion in $u1\ s1\ x = (\text{emb } o\ u2\ s2)\ x$ **for** $\text{emb } u1\ u2\ s1\ s2\ x \Rightarrow$
 $\langle \text{coinduction arbitrary: } x, \text{fastforce} \rangle$

lemma *emb-commute*:

$\wedge s. \text{unfold-lts } (\text{lts-of-dlts-emb } o\ s) = \text{lts-of-dlts } o\ \text{unfold-dlts } s$
 $\wedge s. \text{unfold-gen } (\text{gen-of-lmc-emb } o\ s) = \text{gen-of-lmc } o\ \text{unfold-lmc } s$
 $\wedge s. \text{unfold-lmdp } (\text{lmdp-of-lmc-emb } o\ s) = \text{lmdp-of-lmc } o\ \text{unfold-lmc } s$
 $\wedge s. \text{unfold-react } (\text{react-of-dlts-emb } o\ s) = \text{react-of-dlts } o\ \text{unfold-dlts } s$
 $\wedge s. \text{unfold-sseg } (\text{sseg-of-lmdp-emb } o\ s) = \text{sseg-of-lmdp } o\ \text{unfold-lmdp } s$
 $\wedge s. \text{unfold-sseg } (\text{sseg-of-lts-emb } o\ s) = \text{sseg-of-lts } o\ \text{unfold-lts } s$
 $\wedge s. \text{unfold-sseg } (\text{ssegopt-of-alt-emb } o\ s) = \text{ssegopt-of-alt } o\ \text{unfold-alt } s$
 $\wedge s. \text{unfold-sseg } (\text{sseg-of-react-emb } o\ s) = \text{sseg-of-react } o\ \text{unfold-react } s$
 $\wedge s. \text{unfold-seg } (\text{seg-of-sseg-emb } o\ s) = \text{seg-of-sseg } o\ \text{unfold-sseg } s$
 $\wedge s. \text{unfold-seg } (\text{seg-of-gen-emb } o\ s) = \text{seg-of-gen } o\ \text{unfold-gen } s$
 $\wedge s. \text{unfold-bun } (\text{bun-of-lts-emb } o\ s) = \text{bun-of-lts } o\ \text{unfold-lts } s$
 $\wedge s. \text{unfold-bun } (\text{bunopt-of-alt-emb } o\ s) = \text{bunopt-of-alt } o\ \text{unfold-alt } s$
 $\wedge s. \text{unfold-bun } (\text{bun-of-gen-emb } o\ s) = \text{bun-of-gen } o\ \text{unfold-gen } s$
 $\wedge s. \text{unfold-pz } (\text{pz-of-seg-emb } o\ s) = \text{pz-of-seg } o\ \text{unfold-seg } s$
 $\wedge s. \text{unfold-pz } (\text{pz-of-bun-emb } o\ s) = \text{pz-of-bun } o\ \text{unfold-bun } s$
 $\wedge s. \text{unfold-str } (\text{str-of-mc-emb } o\ s) = \text{str-of-mc } o\ \text{unfold-mc } s$
 $\wedge s. \text{unfold-alt } (\text{alt-of-str-emb } o\ s) = \text{alt-of-str } o\ \text{unfold-str } s$
 $\wedge s. \text{unfold-sseg } (\text{ssegopt-of-sseg-emb } o\ s) = \text{ssegopt-of-sseg } o\ \text{unfold-sseg } s$
 $\wedge s. \text{unfold-seg } (\text{segopt-of-seg-emb } o\ s) = \text{segopt-of-seg } o\ \text{unfold-seg } s$
 $\wedge s. \text{unfold-bun } (\text{bunopt-of-bun-emb } o\ s) = \text{bunopt-of-bun } o\ \text{unfold-bun } s$
 $\wedge s. \text{unfold-pz } (\text{pzopt-of-pz-emb } o\ s) = \text{pzopt-of-pz } o\ \text{unfold-pz } s$
 $\wedge s. \text{unfold-pz } (\text{pzopt-of-mg-emb } o\ s) = \text{pzopt-of-mg } o\ \text{unfold-mg } s$
 $\wedge s. \text{unfold-mg } (\text{mg-of-pzopt-emb } o\ s) = \text{mg-of-pzopt } o\ \text{unfold-pz } s$
 $\wedge s. \text{unfold-mg } (\text{mg-of-pz-emb } o\ s) = \text{mg-of-pz } o\ \text{unfold-pz } s$
 $\wedge s. \text{unfold-mg } (\text{mg-of-alt1-emb } o\ s) = \text{mg-of-alt1 } o\ \text{unfold-alt } s$
 $\wedge s. \text{unfold-mg } (\text{mg-of-alt2-emb } o\ s) = \text{mg-of-alt2 } o\ \text{unfold-alt } s$
 $\wedge s. \text{unfold-pz } (\text{pz-of-alt1-emb } o\ s) = \text{pz-of-alt1 } o\ \text{unfold-alt } s$
 $\wedge s. \text{unfold-pz } (\text{pz-of-alt2-emb } o\ s) = \text{pz-of-alt2 } o\ \text{unfold-alt } s$
by *commute-prover+*

method *inj-prover* =

intro injI,

match conclusion in $x = y$ **for** $x\ y \Rightarrow \langle \text{coinduction arbitrary: } x\ y, \text{fastforce} \rangle$

lemma *inj*:

inj lts-of-dlts
inj react-of-dlts
inj gen-of-lmc
inj lmdp-of-lmc
inj sseg-of-lmdp
inj sseg-of-react
inj sseg-of-lts
inj ssegopt-of-alt
inj seg-of-gen
inj seg-of-sseg
inj bun-of-lts
inj bunopt-of-alt
inj bun-of-gen
inj pz-of-seg
inj pz-of-bun
inj str-of-mc
inj alt-of-str
inj ssegopt-of-sseg
inj segopt-of-seg
inj bunopt-of-bun
inj pzopt-of-pz
inj pzopt-of-mg
inj mg-of-pzopt

inj mg-of-pz
inj mg-of-alt1
inj mg-of-alt2
inj pz-of-alt1
inj pz-of-alt2
by *inj-prover*+

end

lemma *hierarchy*:

$\bigwedge s1\ s2. \text{bisimilar-dlts } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-lts } (\text{lts-of-dlts-emb } o\ s1) (\text{lts-of-dlts-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lmc } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-gen } (\text{gen-of-lmc-emb } o\ s1) (\text{gen-of-lmc-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lmc } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-lmdp } (\text{lmdp-of-lmc-emb } o\ s1) (\text{lmdp-of-lmc-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-dlts } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-react } (\text{react-of-dlts-emb } o\ s1) (\text{react-of-dlts-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lmdp } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{sseg-of-lmdp-emb } o\ s1) (\text{sseg-of-lmdp-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-lts } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{sseg-of-lts-emb } o\ s1) (\text{sseg-of-lts-emb } o\ s2)\ x\ y$
 $\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-sseg } (\text{ssegopt-of-alt-emb } o\ s1) (\text{ssegopt-of-alt-emb } o\ s2)\ x\ y$

$(s\text{segopt-of-alt-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-react } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-sseg } (s\text{seg-of-react-emb } o \ s1)$
 $(s\text{seg-of-react-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-sseg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (s\text{eg-of-sseg-emb } o \ s1) \ (s\text{eg-of-sseg-emb } o \ s2) \ x \ y$
 $\bigwedge (s1 :: - \Rightarrow ('a \ \text{option} \times - \ \text{pmf}) \ \text{set}[-]) \ s2.$
 $\text{bisimilar-sseg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (s\text{eg-of-sseg-emb } o \ s1) \ (s\text{eg-of-sseg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-gen } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (s\text{eg-of-gen-emb } o \ s1) \ (s\text{eg-of-gen-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-lts } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (b\text{un-of-lts-emb } o \ s1) \ (b\text{un-of-lts-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-alt } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (b\text{unopt-of-alt-emb } o \ s1) \ (b\text{unopt-of-alt-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-gen } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (b\text{un-of-gen-emb } o \ s1) \ (b\text{un-of-gen-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-seg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (p\text{z-of-seg-emb } o \ s1) \ (p\text{z-of-seg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-bun } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (p\text{z-of-bun-emb } o \ s1) \ (p\text{z-of-bun-emb } o \ s2) \ x \ y$
 $\bigwedge (s1 :: - \Rightarrow ('a \ \text{option} \times -) \ \text{pmf} \ \text{set}[-]) \ s2.$
 $\text{bisimilar-seg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (p\text{z-of-seg-emb } o \ s1) \ (p\text{z-of-seg-emb } o \ s2) \ x \ y$
 $\bigwedge (s1 :: - \Rightarrow (('a \ \text{option} \times -) \ \text{set}[-]) \ \text{pmf}) \ s2.$
 $\text{bisimilar-bun } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (p\text{z-of-bun-emb } o \ s1) \ (p\text{z-of-bun-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-mc } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-str } (s\text{tr-of-mc-emb } o \ s1) \ (s\text{tr-of-mc-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-sseg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-sseg } (s\text{segopt-of-sseg-emb } o \ s1)$
 $(s\text{segopt-of-sseg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-seg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-seg } (s\text{egopt-of-seg-emb } o \ s1) \ (s\text{egopt-of-seg-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-bun } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-bun } (b\text{unopt-of-bun-emb } o \ s1)$
 $(b\text{unopt-of-bun-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-pz } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (p\text{zopt-of-pz-emb } o \ s1) \ (p\text{zopt-of-pz-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-str } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-alt } (a\text{lt-of-str-emb } o \ s1) \ (a\text{lt-of-str-emb } o \ s2) \ x \ y$
 $\bigwedge s1 \ s2. \text{bisimilar-mg } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-pz } (p\text{zopt-of-mg-emb } o \ s1) \ (p\text{zopt-of-mg-emb } o \ s2) \ x \ y$
unfolding $\text{inj}[THEN \ \text{inj-eq}] \ \text{bisimilar-alt} \ \text{emb-commute} \ o\text{-apply} \ \text{by} \ (\text{rule} \ \text{refl})+$

An edge that would make the graph cyclic

lemma

$\bigwedge s1 \ s2. \text{bisimilar-pz } s1 \ s2 \ x \ y \longleftrightarrow \text{bisimilar-mg } (m\text{g-of-pz-emb } o \ s1) \ (m\text{g-of-pz-emb } o \ s2) \ x \ y$
unfolding $\text{inj}[THEN \ \text{inj-eq}] \ \text{bisimilar-alt} \ \text{emb-commute} \ o\text{-apply} \ \text{by} \ (\text{rule} \ \text{refl})+$

Some redundant (historic) transitive edges

lemma

$\bigwedge s1\ s2. \text{bisimilar-pz } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-mg } (\text{mg-of-pzopt-emb } o\ s1)\ (\text{mg-of-pzopt-emb } o\ s2)\ x\ y$

$\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-mg } (\text{mg-of-alt1-emb } o\ s1)\ (\text{mg-of-alt1-emb } o\ s2)\ x\ y$

$\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-mg } (\text{mg-of-alt2-emb } o\ s1)\ (\text{mg-of-alt2-emb } o\ s2)\ x\ y$

$\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-pz } (\text{pz-of-alt1-emb } o\ s1)\ (\text{pz-of-alt1-emb } o\ s2)\ x\ y$

$\bigwedge s1\ s2. \text{bisimilar-alt } s1\ s2\ x\ y \longleftrightarrow \text{bisimilar-pz } (\text{pz-of-alt2-emb } o\ s1)\ (\text{pz-of-alt2-emb } o\ s2)\ x\ y$

unfolding $\text{inj}[\text{THEN inj-eq}] \text{bisimilar-alt emb-commute o-apply}$ **by** (rule refl)+

7 Some special proofs

Two views on LTS

lemma $\exists f::((a \times s)\ \text{set} \Rightarrow a \Rightarrow s\ \text{set}). \text{bij } f$

by (fastforce simp: bij-def inj-on-def fun-eq-iff image-iff

intro: exI[of - $\lambda S\ a. \{s. (a, s) \in S\}$] exI[of - $\{(a, b). b \in f\ a\}$ for f])

lemma $\exists f::((a \times s)\ \text{set}[(a \times s)\ \text{set}] \Rightarrow a \Rightarrow s\ \text{set}[s\ \text{set}]). \text{bij } f$

by (auto simp: bij-def inj-on-def fun-eq-iff image-iff bset-eq-iff

intro!: exI[of - $\lambda S\ a. \text{bCollect } (\lambda s. \text{bmember } (a, s)\ S)$]

exI[of - $\text{bCollect } (\lambda(a, b). \text{bmember } b\ (f\ a))$ for f])

mc is trivial

lemma *mc-unit*:

fixes $x\ y :: mc$

shows $x = y$

by (coinduction arbitrary: $x\ y$)

(auto simp: pmf.in-rel map-fst-pair-pmf map-snd-pair-pmf intro: exI[of - pair-pmf $x\ y$ for $x\ y$])

lemma *bisimilar-mc* $s1\ s2\ x\ y$

unfolding *bisimilar-alt* **by** (rule *mc-unit*)

8 Printing the Hierarchy Graph

ML \langle

local

$\text{val trim} = \text{filter } (\text{fn } s \Rightarrow s \langle \rangle \ \text{andalso } s \langle \rangle \ \text{set});$

$\text{fun str-of-T } (\text{Type } (c, Ts)) =$

$\text{implode-space } (\text{trim } [\text{commas } (\text{trim } (\text{map str-of-T } Ts)), \text{Long-Name.base-name}$

$c])$

$| \text{str-of-T } - = ;$

```

fun get-edge thm = thm
  |> Thm.concl-of
  |> HOLogic.dest-Trueprop
  |> HOLogic.dest-eq |> fst
  |> dest-comb |> fst
  |> fastype-of
  |> dest-funT
  |> apply2 str-of-T;

val edges = map get-edge @{thms hierarchy[unfolded bisimilar-alt emb-commute
o-apply, THEN iffD1]};

val nodes = distinct (op =) (maps (fn (x, y) => [x, y]) edges);

val node-graph = map (fn s => ((s, Graph-Display.content-node s []), [] : string
list)) nodes;

val graph = fold (fn (x, y) => fn g =>
  AList.map-entry (fn (x, (y, -)) => x = y) x (cons y) g) edges node-graph

in

val - = Graph-Display.display-graph graph

end
>

```

9 Vardi Systems

```

context notes [[bnf-internals]]
begin
  datatype ('a, 'b, 'k) var0 = PMF ('a × 'b) pmf | BPS ('a × 'b) set['k]
end

inductive var-eq :: ('a, 'b, 'k) var0 ⇒ ('a, 'b, 'k) var0 ⇒ bool (infixl <~> 65)
where
  var-eq-reflp[intro]: x ~ x
| [intro]: PMF (return-pmf (a, x)) ~ BPS (bsingleton (a, x))
| [intro]: BPS (bsingleton (a, x)) ~ PMF (return-pmf (a, x))

lemma var-eq-symp: x ~ y ⇒ y ~ x
  by (auto elim: var-eq.cases)

lemma var-eq-transp: x ~ y ⇒ y ~ z ⇒ x ~ z
  by (auto elim!: var-eq.cases)

quotient-type ('a, 'b, 'k) var = ('a, 'b, 'k) var0 / (~)

```

by (*auto intro: equivpI reflpI sympI transpI var-eq-symp var-eq-transp*)

lift-definition *map-var* :: ('a \Rightarrow 'c) \Rightarrow ('b \Rightarrow 'd) \Rightarrow ('a, 'b, 'k) *var* \Rightarrow ('c, 'd, 'k) *var*
is *map-var0*
by (*auto elim!: var-eq.cases simp: map-bset-bsingleton*)

lift-definition *set1-var* :: ('a, 'b, 'k) *var* \Rightarrow 'a *set*
is *set1-var0*
by (*auto elim!: var-eq.cases*)

lift-definition *set2-var* :: ('a, 'b, 'k) *var* \Rightarrow 'b *set*
is *set2-var0*
by (*auto elim!: var-eq.cases*)

inductive *rel-var* :: ('a \Rightarrow 'c \Rightarrow bool) \Rightarrow ('b \Rightarrow 'd \Rightarrow bool) \Rightarrow ('a, 'b, 'k) *var* \Rightarrow ('c, 'd, 'k) *var* \Rightarrow bool **for** *R S* **where**
set1-var $x \subseteq \{(x, y). R\ x\ y\} \Longrightarrow$ *set2-var* $x \subseteq \{(x, y). S\ x\ y\} \Longrightarrow$
rel-var *R S* (*map-var* *fst* *fst* *x*) (*map-var* *snd* *snd* *x*)

abbreviation (*input*) *var0-of-gen-emb* \equiv *case-option* (*BPS* *bempty*) *PMF*
abbreviation (*input*) *var0-of-lts-emb* \equiv *BPS*

lift-definition *var-of-gen-emb* :: ('a \times 'b) *pmf option* \Rightarrow ('a, 'b, 'k) *var* **is** *var0-of-gen-emb*
.

lift-definition *var-of-lts-emb* :: ('a \times 'b) *set['k]* \Rightarrow ('a, 'b, 'k) *var* **is** *var0-of-lts-emb*
.

abbreviation *bisimilar-var* \equiv *bisimilar* ($\lambda R. rel-var\ (=)\ R$)

lemma *map-var0-eq-BPS-iff[simp]*:
map-var0 *f g z* = *BPS* *X* \longleftrightarrow ($\exists Y. z = BPS\ Y \wedge map-bset\ (map-prod\ f\ g)\ Y = X$)
by (*cases* *z*) *auto*

lemma *map-var0-eq-PMF-iff[simp]*:
map-var0 *f g z* = *PMF* *p* \longleftrightarrow ($\exists q. z = PMF\ q \wedge map-pmf\ (map-prod\ f\ g)\ q = p$)
by (*cases* *z*) *auto*

lemma *bisimilar-lts* *s1 s2 x y* \longleftrightarrow *bisimilar-var* (*var-of-lts-emb* *o* *s1*) (*var-of-lts-emb* *o* *s2*) *x y*
(**is** - \longleftrightarrow *bisimilar-var* (*?emb1* *o* -) (*?emb2* *o* -) - -)

unfolding *bisimilar-def* *o-apply* **proof** *safe*
fix *R* **assume** *R x y* **and**
bis: $\forall x\ y. R\ x\ y \longrightarrow rel-bset\ (rel-prod\ (=)\ R)\ (s1\ x)\ (s2\ y)$
from $\langle R\ x\ y \rangle$ **show** $\exists R. R\ x\ y \wedge (\forall x\ y. R\ x\ y \longrightarrow rel-var\ (=)\ R\ (?emb1\ (s1\ x))\ (?emb2\ (s2\ y)))$
proof (*intro* *exI*[*of* - *R*], *safe*)

```

fix  $x y$ 
assume  $R x y$ 
with  $bis$  have  $*$ :  $rel\text{-}bset (rel\text{-}prod (=) R) (s1 x) (s2 y)$  by  $blast$ 
then obtain  $z$  where
   $set\text{-}bset z \subseteq \{(x, y). rel\text{-}prod (=) R x y\}$   $map\text{-}bset fst z = s1 x map\text{-}bset snd$ 
 $z = s2 y$ 
  by ( $auto simp: bset.in\text{-}rel$ )
then show  $rel\text{-}var (=) R (?emb1 (s1 x)) (?emb2 (s2 y))$ 
unfolding  $rel\text{-}var.simps$  by ( $transfer fixing: z$ )
  ( $force simp: bset.map\text{-}comp o\text{-}def split\text{-}beta prod\text{-}set\text{-}simps$ 
   $intro: exI[of - BPS (map\text{-}bset (\lambda((a,b),(c,d)). ((a,c),(b,d))) z]$ )
qed
next
fix  $R$  assume  $R x y$  and
   $bis: \forall x y. R x y \longrightarrow rel\text{-}var (=) R (?emb1 (s1 x)) (?emb2 (s2 y))$ 
from  $\langle R x y \rangle$  show  $\exists R. R x y \wedge (\forall x y. R x y \longrightarrow rel\text{-}bset (rel\text{-}prod (=) R) (s1$ 
 $x) (s2 y))$ 
proof ( $intro exI[of - R], safe$ )
  fix  $x y$ 
  assume  $R x y$ 
  with  $bis$  have  $rel\text{-}var (=) R (?emb1 (s1 x)) (?emb2 (s2 y))$  by  $blast$ 
  then obtain  $z$  where  $*$ :
     $set1\text{-}var z \subseteq \{(x, y). x = y\}$   $set2\text{-}var z \subseteq \{(x, y). R x y\}$ 
     $?emb1 (s1 x) = map\text{-}var fst fst z ?emb2 (s2 y) = map\text{-}var snd snd z$ 
    by ( $auto simp: rel\text{-}var.simps$ )
  then show  $rel\text{-}bset (rel\text{-}prod (=) R) (s1 x) (s2 y)$ 
    by ( $transfer fixing: s1 s2$ ) ( $fastforce simp: bset.in\text{-}rel bset.map\text{-}comp o\text{-}def$ 
     $map\text{-}pmf\text{-}eq\text{-}return\text{-}pmf\text{-}iff$ 
     $split\text{-}beta[abs\text{-}def] map\text{-}prod\text{-}def subset\text{-}eq split\text{-}beta prod\text{-}set\text{-}defs elim!$ :
     $var\text{-}eq.cases$ 
     $intro: exI[of - map\text{-}bset (\lambda((a,b),(c,d)). ((a,c),(b,d))) z \text{ for } z]$ 
     $exI[of - bsingleton ((a,c),(b,d)) \text{ for } a b c d]$ )
qed
qed

lemma  $bisimilar\text{-}gen s1 s2 x y \longleftrightarrow bisimilar\text{-}var (var\text{-}of\text{-}gen\text{-}emb o s1) (var\text{-}of\text{-}gen\text{-}emb$ 
 $o s2) x y$ 
  ( $is - \longleftrightarrow bisimilar\text{-}var (?emb1 o -) (?emb2 o -) - -$ )
unfolding  $bisimilar\text{-}def o\text{-}apply$  proof  $safe$ 
fix  $R$  assume  $R x y$  and
   $bis: \forall x y. R x y \longrightarrow rel\text{-}option (rel\text{-}pmf (rel\text{-}prod (=) R)) (s1 x) (s2 y)$ 
from  $\langle R x y \rangle$  show  $\exists R. R x y \wedge (\forall x y. R x y \longrightarrow rel\text{-}var (=) R (?emb1 (s1 x))$ 
 $(?emb2 (s2 y)))$ 
proof ( $intro exI[of - R], safe$ )
  fix  $x y$ 
  assume  $R x y$ 
  with  $bis$  have  $*$ :  $rel\text{-}option (rel\text{-}pmf (rel\text{-}prod (=) R)) (s1 x) (s2 y)$  by  $blast$ 
  then show  $rel\text{-}var (=) R (?emb1 (s1 x)) (?emb2 (s2 y))$ 
proof ( $cases s1 x s2 y$   $rule: option.exhaust[case\text{-}product option.exhaust]$ )

```

```

    case None-None then show ?thesis unfolding rel-var.simps
      by (transfer fixing: s1 s2) (auto simp: bempty.rep-eq intro!: exI[of - BPS
bempty])
  next
    case (Some-Some p q)
    with * obtain z where
      set-pmf z  $\subseteq$   $\{(x, y). \text{rel-prod } (=) R x y\}$  map-pmf fst z = p map-pmf snd z
= q
      by (auto simp: pmf.in-rel)
    with Some-Some show ?thesis unfolding rel-var.simps
      by (transfer fixing: s1 s2 z) (force simp: pmf.map-comp o-def split-beta
prod-set-simps
      intro: exI[of - PMF (map-pmf  $\lambda((a,b),(c,d)). ((a,c),(b,d))$ ) z])
  qed auto
  qed
next
fix R assume R x y and
  bis:  $\forall x y. R x y \longrightarrow \text{rel-var } (=) R (?emb1 (s1 x)) (?emb2 (s2 y))$ 
from  $\langle R x y \rangle$  show  $\exists R. R x y \wedge (\forall x y. R x y \longrightarrow$ 
  rel-option (rel-pmf (rel-prod (=) R)) (s1 x) (s2 y))
proof (intro exI[of - R], safe)
  fix x y
  assume R x y
  with bis have rel-var (=) R (?emb1 (s1 x)) (?emb2 (s2 y)) by blast
  then obtain z where *:
    set1-var z  $\subseteq$   $\{(x, y). x = y\}$  set2-var z  $\subseteq$   $\{(x, y). R x y\}$ 
    ?emb1 (s1 x) = map-var fst fst z ?emb2 (s2 y) = map-var snd snd z
  by (auto simp: rel-var.simps)
  then show rel-option (rel-pmf (rel-prod (=) R)) (s1 x) (s2 y)
proof (cases s1 x s2 y rule: option.exhaust[case-product option.exhaust])
  case Some-None
  with * show ?thesis by transfer (auto simp: bempty.rep-eq elim!: var-eq.cases)
next
  case None-Some
  with * show ?thesis by transfer (auto elim!: var-eq.cases)
next
  case (Some-Some p q)
  with * show ?thesis
  by transfer (fastforce simp: subset-eq split-beta prod-set-defs
    elim!: var-eq.cases intro!: rel-pmf-reflI)
  qed simp
  qed
qed

```