

Probabilistic Noninterference

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Abstract

We formalize a probabilistic noninterference for a multi-threaded language with uniform scheduling, where probabilistic behaviour comes from both the scheduler and the individual threads. We define notions probabilistic noninterference in two variants: resumption-based and trace-based. For the resumption-based notions, we prove compositionality w.r.t. the language constructs and establish sound type-system-like syntactic criteria.

This is a formalization of the mathematical development presented in the papers [1, 2]. It is the probabilistic variant of the [Possibilistic Noninterference AFP](#) entry.

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1 Simple While Language with probabilistic choice and parallel execution

```
theory Language-Semantics
imports Interface
begin
```

1.1 Preliminaries

Trivia

```
declare zero-le-mult-iff[simp]
declare split-mult-pos-le[simp]
declare zero-le-divide-iff[simp]
```

```
lemma in-minus-Un[simp]:
assumes  $i \in I$ 
shows  $I - \{i\} \cup \{i\} = I$  and  $\{i\} \cup (I - \{i\}) = I$ 
<proof>
```

```
lemma less-plus-cases[case-names Left Right]:
assumes
*:  $(i::nat) < n1 \implies phi$  and
**:  $\bigwedge i2. i = n1 + i2 \implies phi$ 
shows  $phi$ 
<proof>
```

```
lemma less-plus-elim[elim!, consumes 1, case-names Left Right]:
assumes  $i: (i::nat) < n1 + n2$  and
*:  $i < n1 \implies phi$  and
**:  $\bigwedge i2. [i2 < n2; i = n1 + i2] \implies phi$ 
shows  $phi$ 
<proof>
```

lemma *nth-append-singl*[simp]:
 $i < \text{length } al \implies (al @ [a]) ! i = al ! i$
(proof)

lemma *take-append-singl*[simp]:
assumes $n < \text{length } al$ **shows** $\text{take } n (al @ [a]) = \text{take } n al$
(proof)

lemma *length-unique-prefix*:
 $al1 \leq al \implies al2 \leq al \implies \text{length } al1 = \text{length } al2 \implies al1 = al2$
(proof)

take:

lemma *take-length*[simp]:
 $\text{take } (\text{length } al) al = al$
(proof)

lemma *take-le*:
assumes $n < \text{length } al$
shows $\text{take } n al @ [al ! n] \leq al$
(proof)

lemma *list-less-iff-prefix*: $a < b \iff \text{strict-prefix } a b$
(proof)

lemma *take-lt*:
 $n < \text{length } al \implies \text{take } n al < al$
(proof)

lemma *le-take*:
assumes $n1 \leq n2$
shows $\text{take } n1 al \leq \text{take } n2 al$
(proof)

lemma *inj-take*:
assumes $n1 \leq \text{length } al$ **and** $n2 \leq \text{length } al$
shows $\text{take } n1 al = \text{take } n2 al \iff n1 = n2$
(proof)

lemma *lt-take*: $n1 < n2 \implies n2 \leq \text{length } al \implies \text{take } n1 al < \text{take } n2 al$
(proof)

lsum:

definition *lsum* :: $('a \Rightarrow \text{nat}) \Rightarrow 'a \text{ list} \Rightarrow \text{nat}$ **where**
 $\text{lsum } f al \equiv \text{sum-list } (\text{map } f al)$

lemma *lsum-simps*[simp]:
 $\text{lsum } f [] = 0$

$lsum\ f\ (al\ @\ [a]) = lsum\ f\ al + f\ a$
 $\langle proof \rangle$

lemma *lsum-append*:
 $lsum\ f\ (al\ @\ bl) = lsum\ f\ al + lsum\ f\ bl$
 $\langle proof \rangle$

lemma *lsum-cong[fundef-cong]*:
assumes $\bigwedge a. a \in set\ al \implies f\ a = g\ a$
shows $lsum\ f\ al = lsum\ g\ al$
 $\langle proof \rangle$

lemma *lsum-gt-0[simp]*:
assumes $al \neq []$ **and** $\bigwedge a. a \in set\ al \implies 0 < f\ a$
shows $0 < lsum\ f\ al$
 $\langle proof \rangle$

lemma *lsum-mono[simp]*:
assumes $al \leq bl$
shows $lsum\ f\ al \leq lsum\ f\ bl$
 $\langle proof \rangle$

lemma *lsum-mono2[simp]*:
assumes $f: \bigwedge b. b \in set\ bl \implies f\ b > 0$ **and** $le: al < bl$
shows $lsum\ f\ al < lsum\ f\ bl$
 $\langle proof \rangle$

lemma *lsum-take[simp]*:
 $lsum\ f\ (take\ n\ al) \leq lsum\ f\ al$
 $\langle proof \rangle$

lemma *less-lsum-nchotomy*:
assumes $f: \bigwedge a. a \in set\ al \implies 0 < f\ a$
and $i: (i::nat) < lsum\ f\ al$
shows $\exists n\ j. n < length\ al \wedge j < f\ (al\ !\ n) \wedge i = lsum\ f\ (take\ n\ al) + j$
 $\langle proof \rangle$

lemma *less-lsum-unique*:
assumes $\bigwedge a. a \in set\ al \implies (0::nat) < f\ a$
and $n1 < length\ al \wedge j1 < f\ (al\ !\ n1)$
and $n2 < length\ al \wedge j2 < f\ (al\ !\ n2)$
and $lsum\ f\ (take\ n1\ al) + j1 = lsum\ f\ (take\ n2\ al) + j2$
shows $n1 = n2 \wedge j1 = j2$
 $\langle proof \rangle$

definition *locate-pred* **where**
 $locate\ pred\ f\ al\ (i::nat)\ n\ j \equiv$
 $fst\ n\ j < length\ al \wedge$
 $snd\ n\ j < f\ (al\ !\ (fst\ n\ j)) \wedge$

$i = \text{lsum } f \text{ (take (fst } n\text{-j) al) + (snd } n\text{-j)}$

definition *locate* **where**

$\text{locate } f \text{ al } i \equiv \text{SOME } n\text{-j. locate-pred } f \text{ al } i \text{ } n\text{-j}$

definition *locate1* **where** $\text{locate1 } f \text{ al } i \equiv \text{fst (locate } f \text{ al } i)$

definition *locate2* **where** $\text{locate2 } f \text{ al } i \equiv \text{snd (locate } f \text{ al } i)$

lemma *locate-pred-ex*:

assumes $\bigwedge a. a \in \text{set } al \implies 0 < f a$

and $i < \text{lsum } f \text{ al}$

shows $\exists n\text{-j. locate-pred } f \text{ al } i \text{ } n\text{-j}$

<proof>

lemma *locate-pred-unique*:

assumes $\bigwedge a. a \in \text{set } al \implies 0 < f a$

and $\text{locate-pred } f \text{ al } i \text{ } n1\text{-j1 } \text{locate-pred } f \text{ al } i \text{ } n2\text{-j2}$

shows $n1\text{-j1} = n2\text{-j2}$

<proof>

lemma *locate-locate-pred*:

assumes $\bigwedge a. a \in \text{set } al \implies (0::\text{nat}) < f a$

and $i < \text{lsum } f \text{ al}$

shows $\text{locate-pred } f \text{ al } i \text{ (locate } f \text{ al } i)$

<proof>

lemma *locate-locate-pred-unique*:

assumes $\bigwedge a. a \in \text{set } al \implies (0::\text{nat}) < f a$

and $\text{locate-pred } f \text{ al } i \text{ } n\text{-j}$

shows $n\text{-j} = \text{locate } f \text{ al } i$

<proof>

lemma *locate*:

assumes $\bigwedge a. a \in \text{set } al \implies 0 < f a$

and $i < \text{lsum } f \text{ al}$

shows $\text{locate1 } f \text{ al } i < \text{length } al \wedge$

$\text{locate2 } f \text{ al } i < f \text{ (al ! (locate1 } f \text{ al } i)) \wedge$

$i = \text{lsum } f \text{ (take (locate1 } f \text{ al } i) al) + (\text{locate2 } f \text{ al } i)$

<proof>

lemma *locate-unique*:

assumes $\bigwedge a. a \in \text{set } al \implies 0 < f a$

and $n < \text{length } al$ **and** $j < f \text{ (al ! } n)$ **and** $i = \text{lsum } f \text{ (take } n \text{ al) + } j$

shows $n = \text{locate1 } f \text{ al } i \wedge j = \text{locate2 } f \text{ al } i$

<proof>

sum:

lemma *sum-2[simp]*:

$\text{sum } f \text{ \{..< 2\}} = f 0 + f \text{ (Suc } 0)$

<proof>

lemma *inj-Plus[simp]*:
inj ((+) (a::nat))
<proof>

lemma *inj-on-Plus[simp]*:
inj-on ((+) (a::nat)) A
<proof>

lemma *Plus-int[simp]*:
fixes a :: nat
shows (+) b ‘ {..*a*} = {b ..< b + a}
<proof>

lemma *sum-minus[simp]*:
fixes a :: nat
shows *sum f* {a ..< a + b} = *sum* (%x. *f* (a + x)) {..*b*}
<proof>

lemma *sum-Un-introL*:
assumes A1 = B1 Un C1 **and** x = x1 + x2
finite A1 **and**
B1 Int C1 = {} **and**
sum f1 B1 = x1 **and** *sum f1* C1 = x2
shows *sum f1* A1 = x
<proof>

lemma *sum-Un-intro*:
assumes A1 = B1 Un C1 **and** A2 = B2 Un C2 **and**
finite A1 **and** *finite* A2 **and**
B1 Int C1 = {} **and** B2 Int C2 = {} **and**
sum f1 B1 = *sum f2* B2 **and** *sum f1* C1 = *sum f2* C2
shows *sum f1* A1 = *sum f2* A2
<proof>

lemma *sum-UN-introL*:
assumes A1: A1 = (UN n : N. B1 n) **and** a2: a2 = *sum b2* N **and**
fin: *finite* N \wedge n. n \in N \implies *finite* (B1 n) **and**
int: \bigwedge m n. {m, n} \subseteq N \wedge m \neq n \implies B1 m \cap B1 n = {} **and**
ss: \bigwedge n. n \in N \implies *sum f1* (B1 n) = b2 n
shows *sum f1* A1 = a2 (**is** ?L = a2)
<proof>

lemma *sum-UN-intro*:
assumes A1: A1 = (UN n : N. B1 n) **and** A2: A2 = (UN n : N. B2 n) **and**
fin: *finite* N \wedge n. n \in N \implies *finite* (B1 n) \wedge *finite* (B2 n) **and**
int: \bigwedge m n. {m, n} \subseteq N \wedge m \neq n \implies B1 m \cap B1 n = {} \wedge m n. {m, n} \subseteq N
 \implies B2 m \cap B2 n = {} **and**

ss: $\bigwedge n. n \in N \implies \text{sum } f1 (B1\ n) = \text{sum } f2 (B2\ n)$
shows $\text{sum } f1\ A1 = \text{sum } f2\ A2$ (**is** ?L = ?R)
 <proof>

lemma *sum-Minus-intro*:
fixes $f1 :: 'a1 \Rightarrow \text{real}$ **and** $f2 :: 'a2 \Rightarrow \text{real}$
assumes $B1 = A1 - \{a1\}$ **and** $B2 = A2 - \{a2\}$ **and**
 $a1 : A1$ **and** $a2 : A2$ **and** *finite* $A1$ **and** *finite* $A2$
 $\text{sum } f1\ A1 = \text{sum } f2\ A2$ **and** $f1\ a1 = f2\ a2$
shows $\text{sum } f1\ B1 = \text{sum } f2\ B2$
 <proof>

lemma *sum-singl-intro*:
assumes $b = f\ a$
and *finite* A **and** $a \in A$
and $\bigwedge a'. \llbracket a' \in A; a' \neq a \rrbracket \implies f\ a' = 0$
shows $\text{sum } f\ A = b$
 <proof>

lemma *sum-all0-intro*:
assumes $b = 0$
and $\bigwedge a. a \in A \implies f\ a = 0$
shows $\text{sum } f\ A = b$
 <proof>

lemma *sum-1*:
assumes I : *finite* I **and** *ss*: $\text{sum } f\ I = 1$ **and** $i: i \in I - I1$ **and** $I1: I1 \subseteq I$
and $f: \bigwedge i. i \in I \implies (0::\text{real}) \leq f\ i$
shows $f\ i \leq 1 - \text{sum } f\ I1$
 <proof>

1.2 Syntax

datatype ('test, 'atom, 'choice) *cmd* =
 Done
 | *Atm* 'atom
 | *Seq* ('test, 'atom, 'choice) *cmd* ('test, 'atom, 'choice) *cmd* (\leftarrow ;; \rightarrow [60, 61] 60)
 | *While* 'test ('test, 'atom, 'choice) *cmd*
 | *Ch* 'choice ('test, 'atom, 'choice) *cmd* ('test, 'atom, 'choice) *cmd*
 | *Par* ('test, 'atom, 'choice) *cmd* list
 | *ParT* ('test, 'atom, 'choice) *cmd* list

fun *noWhile* **where**
noWhile Done \longleftrightarrow True
 | *noWhile* (*Atm* atm) \longleftrightarrow True
 | *noWhile* ($c1$;; $c2$) \longleftrightarrow *noWhile* $c1 \wedge$ *noWhile* $c2$
 | *noWhile* (*While* *tst* c) \longleftrightarrow False
 | *noWhile* (*Ch* $ch\ c1\ c2$) \longleftrightarrow *noWhile* $c1 \wedge$ *noWhile* $c2$

| *noWhile* (*Par cl*) $\longleftrightarrow (\forall c \in \text{set } cl. \text{noWhile } c)$
| *noWhile* (*ParT cl*) $\longleftrightarrow (\forall c \in \text{set } cl. \text{noWhile } c)$

fun *finished* **where**

finished Done $\longleftrightarrow \text{True}$
| *finished (Atm atm)* $\longleftrightarrow \text{False}$
| *finished (c1 ;; c2)* $\longleftrightarrow \text{False}$
| *finished (While tst c)* $\longleftrightarrow \text{False}$
| *finished (Ch ch c1 c2)* $\longleftrightarrow \text{False}$
| *finished (Par cl)* $\longleftrightarrow (\forall c \in \text{set } cl. \text{finished } c)$
| *finished (ParT cl)* $\longleftrightarrow (\forall c \in \text{set } cl. \text{finished } c)$

definition *noWhileL* **where**

noWhileL cl $\equiv \forall c \in \text{set } cl. \text{noWhile } c$

lemma *fin-Par-noWhileL[simp]*:

noWhile (Par cl) $\longleftrightarrow \text{noWhileL } cl$
<proof>

lemma *fin-ParT-noWhileL[simp]*:

noWhile (ParT cl) $\longleftrightarrow \text{noWhileL } cl$
<proof>

declare *noWhile.simps(6)* [*simp del*]

declare *noWhile.simps(7)* [*simp del*]

lemma *noWhileL-intro[intro]*:

assumes $\bigwedge c. c \in \text{set } cl \implies \text{noWhile } c$
shows *noWhileL cl*
<proof>

lemma *noWhileL-fin[simp]*:

assumes *noWhileL cl* **and** $c \in \text{set } cl$
shows *noWhile c*
<proof>

lemma *noWhileL-update[simp]*:

assumes *cl: noWhileL cl* **and** $c': \text{noWhile } c'$
shows *noWhileL (cl[n := c'])*
<proof>

definition *finishedL* **where**

finishedL cl $\equiv \forall c \in \text{set } cl. \text{finished } c$

lemma *finished-Par-finishedL[simp]*:

finished (Par cl) $\longleftrightarrow \text{finishedL } cl$
<proof>

lemma *finished-ParT-finishedL*[simp]:
finished (ParT cl) \longleftrightarrow *finishedL* cl
 ⟨proof⟩

declare *finished.simps*(6) [simp del]
declare *finished.simps*(7) [simp del]

lemma *finishedL-intro*[intro]:
assumes $\bigwedge c. c \in \text{set } cl \implies \text{finished } c$
shows *finishedL* cl
 ⟨proof⟩

lemma *finishedL-finished*[simp]:
assumes *finishedL* cl **and** $c \in \text{set } cl$
shows *finished* c
 ⟨proof⟩

lemma *finishedL-update*[simp]:
assumes *finishedL* cl **and** *finished* c'
shows *finishedL* (cl[n := c'])
 ⟨proof⟩

lemma *finished-fin*[simp]:
finished c \implies *noWhile* c
 ⟨proof⟩

lemma *finishedL-noWhileL*[simp]:
finishedL cl \implies *noWhileL* cl
 ⟨proof⟩

locale PL =
fixes
 aval :: 'atom \Rightarrow 'state \Rightarrow 'state **and**
 tval :: 'test \Rightarrow 'state \Rightarrow bool **and**
 cval :: 'choice \Rightarrow 'state \Rightarrow real
assumes
 properCh: $\bigwedge ch\ s. 0 \leq \text{cval } ch\ s \wedge \text{cval } ch\ s \leq 1$
begin

lemma [simp]: $(n::nat) < N \implies 0 \leq 1 / N$ ⟨proof⟩

lemma [simp]: $(n::nat) < N \implies 1 / N \leq 1$ ⟨proof⟩

lemma [simp]: $(n::nat) < N \implies 0 \leq 1 - 1 / N$ ⟨proof⟩

lemma *sum-equal*: $0 < (N::nat) \implies \text{sum } (\lambda n. 1/N) \{..< N\} = 1$
 ⟨proof⟩

fun *proper* **where**

$proper\ Done \longleftrightarrow True$
 $|\ proper\ (Atm\ x) \longleftrightarrow True$
 $| \ proper\ (Seq\ c1\ c2) \longleftrightarrow proper\ c1 \wedge proper\ c2$
 $| \ proper\ (While\ tst\ c) \longleftrightarrow proper\ c$
 $| \ proper\ (Ch\ ch\ c1\ c2) \longleftrightarrow proper\ c1 \wedge proper\ c2$
 $| \ proper\ (Par\ cl) \longleftrightarrow cl \neq [] \wedge (\forall\ c \in set\ cl.\ proper\ c)$
 $| \ proper\ (ParT\ cl) \longleftrightarrow cl \neq [] \wedge (\forall\ c \in set\ cl.\ proper\ c)$

definition *properL* **where**

$properL\ cl \equiv cl \neq [] \wedge (\forall\ c \in set\ cl.\ proper\ c)$

lemma *proper-Par-properL*[simp]:

$proper\ (Par\ cl) \longleftrightarrow properL\ cl$
 $\langle proof \rangle$

lemma *proper-ParT-properL*[simp]:

$proper\ (ParT\ cl) \longleftrightarrow properL\ cl$
 $\langle proof \rangle$

declare *proper.simps*(6) [simp del]

declare *proper.simps*(7) [simp del]

lemma *properL-intro*[intro]:

$\llbracket cl \neq []; \bigwedge c. c \in set\ cl \implies proper\ c \rrbracket \implies properL\ cl$
 $\langle proof \rangle$

lemma *properL-notEmp*[simp]: $properL\ cl \implies cl \neq []$

$\langle proof \rangle$

lemma *properL-proper*[simp]:

$\llbracket properL\ cl; c \in set\ cl \rrbracket \implies proper\ c$
 $\langle proof \rangle$

lemma *properL-update*[simp]:

assumes $cl: properL\ cl$ **and** $c': proper\ c'$

shows $properL\ (cl[n := c'])$

$\langle proof \rangle$

lemma *proper-induct*[consumes 1, case-names Done Atm Seq While Ch Par ParT]:

assumes $*$: $proper\ c$

and *Done*: $phi\ Done$

and *Atm*: $\bigwedge atm. phi\ (Atm\ atm)$

and *Seq*: $\bigwedge c1\ c2. \llbracket phi\ c1; phi\ c2 \rrbracket \implies phi\ (c1\ ;;\ c2)$

and *While*: $\bigwedge tst\ c. phi\ c \implies phi\ (While\ tst\ c)$

and *Ch*: $\bigwedge ch\ c1\ c2. \llbracket phi\ c1; phi\ c2 \rrbracket \implies phi\ (Ch\ ch\ c1\ c2)$

and *Par*: $\bigwedge cl. \llbracket properL\ cl; \bigwedge c. c \in set\ cl \implies phi\ c \rrbracket \implies phi\ (Par\ cl)$

and *ParT*: $\bigwedge cl. \llbracket properL\ cl; \bigwedge c. c \in set\ cl \implies phi\ c \rrbracket \implies phi\ (ParT\ cl)$

shows $phi\ c$

$\langle proof \rangle$

1.2.1 Operational Small-Step Semantics

definition $theFT\ cl \equiv \{n. n < length\ cl \wedge finished\ (cl!n)\}$

definition $theNFT\ cl \equiv \{n. n < length\ cl \wedge \neg finished\ (cl!n)\}$

lemma $finite-theFT[simp]: finite\ (theFT\ cl)$
 $\langle proof \rangle$

lemma $theFT-length[simp]: n \in theFT\ cl \implies n < length\ cl$
 $\langle proof \rangle$

lemma $theFT-finished[simp]: n \in theFT\ cl \implies finished\ (cl!n)$
 $\langle proof \rangle$

lemma $finite-theNFT[simp]: finite\ (theNFT\ cl)$
 $\langle proof \rangle$

lemma $theNFT-length[simp]: n \in theNFT\ cl \implies n < length\ cl$
 $\langle proof \rangle$

lemma $theNFT-notFinished[simp]: n \in theNFT\ cl \implies \neg finished\ (cl!n)$
 $\langle proof \rangle$

lemma $theFT-Int-theNFT[simp]:$
 $theFT\ cl\ Int\ theNFT\ cl = \{\}$ **and** $theNFT\ cl\ Int\ theFT\ cl = \{\}$
 $\langle proof \rangle$

lemma $theFT-Un-theNFT[simp]:$
 $theFT\ cl\ Un\ theNFT\ cl = \{.. < length\ cl\}$ **and**
 $theNFT\ cl\ Un\ theFT\ cl = \{.. < length\ cl\}$
 $\langle proof \rangle$

lemma $in-theFT-theNFT[simp]:$
assumes $n1 \in theFT\ cl$ **and** $n2 \in theNFT\ cl$
shows $n1 \neq n2$ **and** $n2 \neq n1$
 $\langle proof \rangle$

definition $WtFT\ cl \equiv sum\ (\lambda\ (n::nat). 1/(length\ cl))\ (theFT\ cl)$

definition $WtNFT\ cl \equiv sum\ (\lambda\ (n::nat). 1/(length\ cl))\ (theNFT\ cl)$

lemma $WtFT-WtNFT[simp]:$
assumes $0 < length\ cl$
shows $WtFT\ cl + WtNFT\ cl = 1$ (**is** $?A = 1$)
 $\langle proof \rangle$

lemma *WtNFT-1-WtFT*: $0 < \text{length } cl \implies \text{WtNFT } cl = 1 - \text{WtFT } cl$
<proof>

lemma *WtNFT-WtFT-1[simp]*:
assumes $0 < \text{length } cl$ **and** $\text{WtFT } cl \neq 1$
shows $\text{WtNFT } cl / (1 - \text{WtFT } cl) = 1$ (**is** ?A / ?B = 1)
<proof>

lemma *WtFT-ge-0[simp]*: $\text{WtFT } cl \geq 0$
<proof>

lemma *WtFT-le-1[simp]*: $\text{WtFT } cl \leq 1$ (**is** ?L ≤ 1)
<proof>

lemma *le-1-WtFT[simp]*: $0 \leq 1 - \text{WtFT } cl$ (**is** $0 \leq ?R$)
<proof>

lemma *WtFT-lt-1[simp]*: $\text{WtFT } cl \neq 1 \implies \text{WtFT } cl < 1$
<proof>

lemma *lt-1-WtFT[simp]*: $\text{WtFT } cl \neq 1 \implies 0 < 1 - \text{WtFT } cl$
<proof>

lemma *notFinished-WtFT[simp]*:
assumes $n < \text{length } cl$ **and** $\neg \text{finished } (cl ! n)$
shows $1 / \text{length } cl \leq 1 - \text{WtFT } cl$
<proof>

fun *brn* :: ('test, 'atom, 'choice) cmd \Rightarrow nat **where**
| *brn Done* = 1
| *brn (Atm atm)* = 1
| *brn (c1 ;; c2)* = *brn c1*
| *brn (While tst c)* = 1
| *brn (Ch ch c1 c2)* = 2
| *brn (Par cl)* = *lsum brn cl*
| *brn (ParT cl)* = *lsum brn cl*

lemma *brn-gt-0*: *proper c* $\implies 0 < \text{brn } c$
<proof>

lemma *brn-gt-0-L*: $\llbracket \text{properL } cl; c \in \text{set } cl \rrbracket \implies 0 < \text{brn } c$
<proof>

definition *locateT* \equiv *locate1 brn* **definition** *locateI* \equiv *locate2 brn*

definition *brnL cl n* \equiv *lsum brn (take n cl)*

lemma *brnL-lsum*: $\text{brnL } cl \ (\text{length } cl) = \text{lsum } \text{brn } cl$
<proof>

lemma *brnL-unique*:

assumes *properL* cl **and** $n1 < \text{length } cl \wedge j1 < \text{brn } (cl ! n1)$
and $n2 < \text{length } cl \wedge j2 < \text{brn } (cl ! n2)$ **and** $\text{brnL } cl \ n1 + j1 = \text{brnL } cl \ n2 + j2$
shows $n1 = n2 \wedge j1 = j2$
<proof>

lemma *brn-Par-simp[simp]*: $\text{brn } (\text{Par } cl) = \text{brnL } cl \ (\text{length } cl)$
<proof>

lemma *brn-ParT-simp[simp]*: $\text{brn } (\text{ParT } cl) = \text{brnL } cl \ (\text{length } cl)$
<proof>

declare *brn.simps(6)[simp del]* **declare** *brn.simps(7)[simp del]*

lemma *brnL-0[simp]*: $\text{brnL } cl \ 0 = 0$
<proof>

lemma *brnL-Suc[simp]*: $n < \text{length } cl \implies \text{brnL } cl \ (\text{Suc } n) = \text{brnL } cl \ n + \text{brn } (cl ! n)$
<proof>

lemma *brnL-mono[simp]*: $n1 \leq n2 \implies \text{brnL } cl \ n1 \leq \text{brnL } cl \ n2$
<proof>

lemma *brnL-mono2[simp]*:
assumes p : *properL* cl **and** n : $n1 < n2$ **and** l : $n2 \leq \text{length } cl$
shows $\text{brnL } cl \ n1 < \text{brnL } cl \ n2$ (**is** $?L < ?R$)
<proof>

lemma *brn-index[simp]*:
assumes n : $n < \text{length } cl$ **and** i : $i < \text{brn } (cl ! n)$
shows $\text{brnL } cl \ n + i < \text{brnL } cl \ (\text{length } cl)$ (**is** $?L < ?R$)
<proof>

lemma *brnL-gt-0[simp]*: $\llbracket \text{properL } cl; 0 < n \rrbracket \implies 0 < \text{brnL } cl \ n$
<proof>

lemma *locateTI*:

assumes *properL* cl **and** $ii < \text{brn } (\text{Par } cl)$
shows
 $\text{locateT } cl \ ii < \text{length } cl \wedge$
 $\text{locateI } cl \ ii < \text{brn } (cl ! (\text{locateT } cl \ ii)) \wedge$
 $ii = \text{brnL } cl \ (\text{locateT } cl \ ii) + \text{locateI } cl \ ii$
<proof>

lemma *locateTI-unique*:

assumes $\text{properL } cl$ **and** $n < \text{length } cl$
and $i < \text{brn } (cl ! n)$ **and** $ii = \text{brnL } cl \ n + i$
shows $n = \text{locateT } cl \ ii \wedge i = \text{locateI } cl \ ii$
 $\langle \text{proof} \rangle$

definition pickFT-pred **where** $\text{pickFT-pred } cl \ n \equiv n < \text{length } cl \wedge \text{finished } (cl ! n)$

definition pickFT **where** $\text{pickFT } cl \equiv \text{SOME } n. \text{pickFT-pred } cl \ n$

lemma pickFT-pred :

assumes $\text{WtFT } cl = 1$ **shows** $\exists n. \text{pickFT-pred } cl \ n$
 $\langle \text{proof} \rangle$

lemma $\text{pickFT-pred-pickFT}$: $\text{WtFT } cl = 1 \implies \text{pickFT-pred } cl \ (\text{pickFT } cl)$
 $\langle \text{proof} \rangle$

lemma $\text{pickFT-length[simp]}$: $\text{WtFT } cl = 1 \implies \text{pickFT } cl < \text{length } cl$
 $\langle \text{proof} \rangle$

lemma $\text{pickFT-finished[simp]}$: $\text{WtFT } cl = 1 \implies \text{finished } (cl ! (\text{pickFT } cl))$
 $\langle \text{proof} \rangle$

lemma $\text{pickFT-theFT[simp]}$: $\text{WtFT } cl = 1 \implies \text{pickFT } cl \in \text{theFT } cl$
 $\langle \text{proof} \rangle$

fun wt-cont-eff **where**

$\text{wt-cont-eff } \text{Done } s \ i = (1, \text{Done}, s)$
 $|$
 $\text{wt-cont-eff } (\text{Atm } atm) \ s \ i = (1, \text{Done}, \text{aval } atm \ s)$
 $|$
 $\text{wt-cont-eff } (c1 \ ; \ ; \ c2) \ s \ i =$
 $(\text{case } \text{wt-cont-eff } c1 \ s \ i \ \text{of}$
 $(x, c1', s') \Rightarrow$
 $\text{if } \text{finished } c1' \ \text{then } (x, c2, s') \ \text{else } (x, c1' \ ; \ ; \ c2, s'))$
 $|$
 $\text{wt-cont-eff } (\text{While } tst \ c) \ s \ i =$
 $(\text{if } \text{tval } tst \ s$
 $\text{then } (1, c \ ; \ ; \ (\text{While } tst \ c), s)$
 $\text{else } (1, \text{Done}, s))$
 $|$
 $\text{wt-cont-eff } (\text{Ch } ch \ c1 \ c2) \ s \ i =$
 $(\text{if } i = 0 \ \text{then } \text{cval } ch \ s \ \text{else } 1 - \text{cval } ch \ s,$
 $\text{if } i = 0 \ \text{then } c1 \ \text{else } c2,$
 $s)$
 $|$
 $\text{wt-cont-eff } (\text{Par } cl) \ s \ ii =$
 $(\text{if } cl ! (\text{locateT } cl \ ii) \in \text{set } cl \ \text{then}$
 $(\text{case } \text{wt-cont-eff}$

```

      (cl ! (locateT cl ii))
      s
      (locateI cl ii) of
(w, c', s') ⇒
  ((1 / (length cl)) * w,
   Par (cl [(locateT cl ii) := c']),
   s')
else undefined)
|
wt-cont-eff (ParT cl) s ii =
  (if cl ! (locateT cl ii) ∈ set cl
   then
     (case wt-cont-eff
        (cl ! (locateT cl ii))
        s
        (locateI cl ii) of
      (w, c', s') ⇒
        (if WtFT cl = 1
         then (if locateT cl ii = pickFT cl ∧ locateI cl ii = 0
              then 1
              else 0)
         else if finished (cl ! (locateT cl ii))
              then 0
              else (1 / (length cl))
                   / (1 - WtFT cl)
                   * w,
                 ParT (cl [(locateT cl ii) := c']),
                 s'))
     else undefined)

```

definition *wt where* $wt\ c\ s\ i = fst\ (wt\text{-cont-eff}\ c\ s\ i)$

definition *cont where* $cont\ c\ s\ i = fst\ (snd\ (wt\text{-cont-eff}\ c\ s\ i))$

definition *eff where* $eff\ c\ s\ i = snd\ (snd\ (wt\text{-cont-eff}\ c\ s\ i))$

lemma *wt-Done[simp]:* $wt\ Done\ s\ i = 1$
 ⟨proof⟩

lemma *wt-Atm[simp]:* $wt\ (Atm\ atm)\ s\ i = 1$
 ⟨proof⟩

lemma *wt-Seq[simp]:*
 $wt\ (c1\ ;;\ c2)\ s = wt\ c1\ s$
 ⟨proof⟩

lemma *wt-While[simp]:* $wt\ (While\ tst\ c)\ s\ i = 1$
 ⟨proof⟩

lemma *wt-Ch-L[simp]*: $wt (Ch\ ch\ c1\ c2)\ s\ 0 = cval\ ch\ s$
<proof>

lemma *wt-Ch-R[simp]*: $wt (Ch\ ch\ c1\ c2)\ s\ (Suc\ n) = 1 - cval\ ch\ s$
<proof>

lemma *cont-Done[simp]*: $cont\ Done\ s\ i = Done$
<proof>

lemma *cont-Atm[simp]*: $cont\ (Atm\ atm)\ s\ i = Done$
<proof>

lemma *cont-Seq-finished[simp]*: $finished\ (cont\ c1\ s\ i) \implies cont\ (c1\ ;;\ c2)\ s\ i = c2$
<proof>

lemma *cont-Seq-notFinished[simp]*:
assumes $\neg\ finished\ (cont\ c1\ s\ i)$
shows $cont\ (c1\ ;;\ c2)\ s\ i = (cont\ c1\ s\ i) ;;\ c2$
<proof>

lemma *cont-Seq-not-eq-finished[simp]*: $\neg\ finished\ c2 \implies \neg\ finished\ (cont\ (Seq\ c1\ c2)\ s\ i)$
<proof>

lemma *cont-While-False[simp]*: $tval\ tst\ s = False \implies cont\ (While\ tst\ c)\ s\ i = Done$
<proof>

lemma *cont-While-True[simp]*: $tval\ tst\ s = True \implies cont\ (While\ tst\ c)\ s\ i = c ;;\ (While\ tst\ c)$
<proof>

lemma *cont-Ch-L[simp]*: $cont\ (Ch\ ch\ c1\ c2)\ s\ 0 = c1$
<proof>

lemma *cont-Ch-R[simp]*: $cont\ (Ch\ ch\ c1\ c2)\ s\ (Suc\ n) = c2$
<proof>

lemma *eff-Done[simp]*: $eff\ Done\ s\ i = s$
<proof>

lemma *eff-Atm[simp]*: $eff\ (Atm\ atm)\ s\ i = aval\ atm\ s$
<proof>

lemma *eff-Seq[simp]*: $eff\ (c1\ ;;\ c2)\ s = eff\ c1\ s$
<proof>

lemma *eff-While[simp]*: $\text{eff } (\text{While } \text{tst } c) s i = s$
 $\langle \text{proof} \rangle$

lemma *eff-Ch[simp]*: $\text{eff } (\text{Ch } \text{ch } c1 c2) s i = s$
 $\langle \text{proof} \rangle$

lemma *brnL-nchotomy*:

assumes *properL cl and ii < brnL cl (length cl)*

shows $\exists n i. n < \text{length } cl \wedge i < \text{brn } (cl ! n) \wedge ii = \text{brnL } cl n + i$
 $\langle \text{proof} \rangle$

corollary *brnL-cases[consumes 2, case-names Local, elim]*:

assumes *properL cl and ii < brnL cl (length cl) and*

$\bigwedge n i. \llbracket n < \text{length } cl; i < \text{brn } (cl ! n); ii = \text{brnL } cl n + i \rrbracket \implies phi$

shows *phi*

$\langle \text{proof} \rangle$

lemma *wt-cont-eff-Par[simp]*:

assumes *p: properL cl*

and *n: n < length cl and i: i < brn (cl ! n)*

shows

$\text{wt } (\text{Par } cl) s (\text{brnL } cl n + i) =$
 $1 / (\text{length } cl) * \text{wt } (cl ! n) s i$
 $(\text{is } ?wL = ?wR)$

$\text{cont } (\text{Par } cl) s (\text{brnL } cl n + i) =$
 $\text{Par } (cl [n := \text{cont } (cl ! n) s i])$
 $(\text{is } ?mL = ?mR)$

$\text{eff } (\text{Par } cl) s (\text{brnL } cl n + i) =$
 $\text{eff } (cl ! n) s i$
 $(\text{is } ?eL = ?eR)$
 $\langle \text{proof} \rangle$

lemma *cont-eff-ParT[simp]*:

assumes *p: properL cl*

and *n: n < length cl and i: i < brn (cl ! n)*

shows

$\text{cont } (\text{ParT } cl) s (\text{brnL } cl n + i) =$
 $\text{ParT } (cl [n := \text{cont } (cl ! n) s i])$
 $(\text{is } ?mL = ?mR)$

$\text{eff } (\text{ParT } cl) s (\text{brnL } cl n + i) =$
 $\text{eff } (cl ! n) s i$
 $(\text{is } ?eL = ?eR)$
 $\langle \text{proof} \rangle$

lemma *wt-ParT-WtFT-pickFT-0[simp]*:
assumes p : *properL cl* **and** $WtFT$: $WtFT\ cl = 1$
shows $wt\ (ParT\ cl)\ s\ (brnL\ cl\ (pickFT\ cl)) = 1$
(is $?wL = 1$)
 $\langle proof \rangle$

lemma *wt-ParT-WtFT-notPickFT-0[simp]*:
assumes p : *properL cl* **and** n : $n < length\ cl$ **and** i : $i < brn\ (cl!\ n)$
and $WtFT$: $WtFT\ cl = 1$ **and** ni : $n = pickFT\ cl \longrightarrow i \neq 0$
shows $wt\ (ParT\ cl)\ s\ (brnL\ cl\ n + i) = 0$ **(is** $?wL = 0$)
 $\langle proof \rangle$

lemma *wt-ParT-notWtFT-finished[simp]*:
assumes p : *properL cl* **and** n : $n < length\ cl$ **and** i : $i < brn\ (cl!\ n)$
and $WtFT$: $WtFT\ cl \neq 1$ **and** f : *finished (cl! n)*
shows $wt\ (ParT\ cl)\ s\ (brnL\ cl\ n + i) = 0$ **(is** $?wL = 0$)
 $\langle proof \rangle$

lemma *wt-cont-eff-ParT-notWtFT-notFinished[simp]*:
assumes p : *properL cl* **and** n : $n < length\ cl$ **and** i : $i < brn\ (cl!\ n)$
and $WtFT$: $WtFT\ cl \neq 1$ **and** nf : $\neg finished\ (cl!\ n)$
shows $wt\ (ParT\ cl)\ s\ (brnL\ cl\ n + i) =$
 $(1 / (length\ cl)) / (1 - WtFT\ cl) * wt\ (cl!\ n)\ s\ i$ **(is** $?wL = ?wR$)
 $\langle proof \rangle$

lemma *wt-ge-0[simp]*:
assumes *proper c* **and** $i < brn\ c$
shows $0 \leq wt\ c\ s\ i$
 $\langle proof \rangle$

lemma *wt-le-1[simp]*:
assumes *proper c* **and** $i < brn\ c$
shows $wt\ c\ s\ i \leq 1$
 $\langle proof \rangle$

abbreviation *fromPlus* ($\langle (1\ \{-..\langle +\ -\}) \rangle$) **where**
 $\{a\ ..\langle +\ b\ \} \equiv \{a\ ..\langle a + b\ \}$

lemma *brnL-UN*:
assumes *properL cl*
shows $\{..\langle brnL\ cl\ (length\ cl)\ \} = (\bigcup\ n < length\ cl.\ \{brnL\ cl\ n\ ..\langle +\ brn\ (cl!\ n)\ \})$
(is $?L = (\bigcup\ n < length\ cl.\ ?R\ n)$)
 $\langle proof \rangle$

lemma *brnL-Int-It*:
assumes $n12$: $n1 < n2$ **and** $n2$: $n2 < length\ cl$
shows
 $\{brnL\ cl\ n1\ ..\langle +\ brn\ (cl!\ n1)\ \} \cap \{brnL\ cl\ n2\ ..\langle +\ brn\ (cl!\ n2)\ \} = \{\}$

$\langle proof \rangle$

lemma *brnL-Int*:

assumes $n1 \neq n2$ **and** $n1 < length\ cl$ **and** $n2 < length\ cl$

shows $\{brnL\ cl\ n1\ ..<+ \ brn\ (cl!n1)\} \cap \{brnL\ cl\ n2\ ..<+ \ brn\ (cl!n2)\} = \{\}$

$\langle proof \rangle$

lemma *sum-wt-Par-sub[simp]*:

assumes $cl: properL\ cl$ **and** $n: n < length\ cl$ **and** $I: I \subseteq \{..< \ brn\ (cl!\ n)\}$

shows $sum\ (wt\ (Par\ cl)\ s)\ ((+)\ (brnL\ cl\ n)\ 'I) =$

$$1 / (length\ cl) * sum\ (wt\ (cl!\ n)\ s)\ I\ (\mathbf{is}\ ?L = ?wSch * ?R)$$

$\langle proof \rangle$

lemma *sum-wt-Par[simp]*:

assumes $cl: properL\ cl$ **and** $n: n < length\ cl$

shows $sum\ (wt\ (Par\ cl)\ s)\ \{brnL\ cl\ n\ ..<+ \ brn\ (cl!n)\} =$

$$1 / (length\ cl) * sum\ (wt\ (cl!\ n)\ s)\ \{..< \ brn\ (cl!\ n)\}\ (\mathbf{is}\ ?L = ?W * ?R)$$

$\langle proof \rangle$

lemma *sum-wt-ParT-sub-WtFT-pickFT-0[simp]*:

assumes $cl: properL\ cl$ **and** $nf: WtFT\ cl = 1$

and $I: I \subseteq \{..< \ brn\ (cl!\ (pickFT\ cl))\}\ 0 \in I$

shows $sum\ (wt\ (ParT\ cl)\ s)\ ((+)\ (brnL\ cl\ (pickFT\ cl))\ 'I) = 1\ (\mathbf{is}\ ?L = 1)$

$\langle proof \rangle$

lemma *sum-wt-ParT-sub-WtFT-pickFT-0-2[simp]*:

assumes $cl: properL\ cl$ **and** $nf: WtFT\ cl = 1$

and $II: II \subseteq \{..< \ brnL\ cl\ (length\ cl)\}\ brnL\ cl\ (pickFT\ cl) \in II$

shows $sum\ (wt\ (ParT\ cl)\ s)\ II = 1\ (\mathbf{is}\ ?L = 1)$

$\langle proof \rangle$

lemma *sum-wt-ParT-sub-WtFT-notPickFT-0[simp]*:

assumes $cl: properL\ cl$ **and** $nf: WtFT\ cl = 1$ **and** $n: n < length\ cl$

and $I: I \subseteq \{..< \ brn\ (cl!\ n)\}$ **and** $nI: n = pickFT\ cl \longrightarrow 0 \notin I$

shows $sum\ (wt\ (ParT\ cl)\ s)\ ((+)\ (brnL\ cl\ n)\ 'I) = 0\ (\mathbf{is}\ ?L = 0)$

$\langle proof \rangle$

lemma *sum-wt-ParT-sub-notWtFT-finished[simp]*:

assumes $cl: properL\ cl$ **and** $nf: WtFT\ cl \neq 1$

and $n: n < length\ cl$ **and** $chn: finished\ (cl!n)$ **and** $I: I \subseteq \{..< \ brn\ (cl!\ n)\}$

shows $sum\ (wt\ (ParT\ cl)\ s)\ ((+)\ (brnL\ cl\ n)\ 'I) = 0\ (\mathbf{is}\ ?L = 0)$

$\langle proof \rangle$

lemma *sum-wt-ParT-sub-notWtFT-notFinished[simp]*:

assumes $cl: properL\ cl$ **and** $nf: WtFT\ cl \neq 1$ **and** $n: n < length\ cl$

and $chn: \neg finished\ (cl!n)$ **and** $I: I \subseteq \{..< \ brn\ (cl!\ n)\}$

shows

$sum\ (wt\ (ParT\ cl)\ s)\ ((+)\ (brnL\ cl\ n)\ 'I) =$

$$(1 / (length\ cl)) / (1 - WtFT\ cl) * sum\ (wt\ (cl!\ n)\ s)\ I$$

(is ?L = ?w / (1 - ?wF) * ?R)
 ⟨proof⟩

lemma *sum-wt-ParT-WtFT-pickFT-0[simp]*:
assumes *cl: properL cl and nf: WtFT cl = 1*
shows $\text{sum } (wt \text{ (ParT } cl) \text{ s}) \{brnL \text{ cl } (pickFT \text{ cl}) ..<+ brn \text{ (cl ! (pickFT cl))}\} = 1$
 ⟨proof⟩

lemma *sum-wt-ParT-WtFT-notPickFT-0[simp]*:
assumes *cl: properL cl and nf: WtFT cl = 1 and n: n < length cl n ≠ pickFT cl*
shows $\text{sum } (wt \text{ (ParT } cl) \text{ s}) \{brnL \text{ cl } n ..<+ brn \text{ (cl!n)}\} = 0$
 ⟨proof⟩

lemma *sum-wt-ParT-notWtFT-finished[simp]*:
assumes *cl: properL cl and WtFT cl ≠ 1*
and *n: n < length cl and cln: finished (cl!n)*
shows $\text{sum } (wt \text{ (ParT } cl) \text{ s}) \{brnL \text{ cl } n ..<+ brn \text{ (cl!n)}\} = 0$
 ⟨proof⟩

lemma *sum-wt-ParT-notWtFT-notFinished[simp]*:
assumes *cl: properL cl and nf: WtFT cl ≠ 1*
and *n: n < length cl and cln: ¬ finished (cl!n)*
shows
 $\text{sum } (wt \text{ (ParT } cl) \text{ s}) \{brnL \text{ cl } n ..<+ brn \text{ (cl!n)}\} =$
 $(1 / (\text{length } cl)) / (1 - WtFT \text{ cl}) *$
 $\text{sum } (wt \text{ (cl ! } n) \text{ s}) \{..< brn \text{ (cl ! } n)\}$
 ⟨proof⟩

lemma *sum-wt[simp]*:
assumes *proper c*
shows $\text{sum } (wt \text{ c s}) \{..< brn \text{ c}\} = 1$
 ⟨proof⟩

lemma *proper-cont[simp]*:
assumes *proper c and i < brn c*
shows *proper (cont c s i)*
 ⟨proof⟩

lemma *sum-subset-le-1[simp]*:
assumes **: proper c and **: I ⊆ {..< brn c}*
shows $\text{sum } (wt \text{ c s}) \text{ I} \leq 1$
 ⟨proof⟩

lemma *sum-le-1[simp]*:
assumes **: proper c and **: i < brn c*
shows $\text{sum } (wt \text{ c s}) \{..i\} \leq 1$
 ⟨proof⟩

1.2.2 Operations on configurations

definition $cont\text{-}eff\ cf\ b = snd\ (wt\text{-}cont\text{-}eff\ (fst\ cf)\ (snd\ cf)\ b)$

lemma $cont\text{-}eff$: $cont\text{-}eff\ cf\ b = (cont\ (fst\ cf)\ (snd\ cf)\ b, eff\ (fst\ cf)\ (snd\ cf)\ b)$
 $\langle proof \rangle$

end

end

2 Resumption-Based Noninterference

theory *Resumption-Based*
imports *Language-Semantics*
begin

type-synonym $'a\ rel = ('a \times 'a)\ set$

2.1 Preliminaries

lemma $int\text{-}emp[simp]$:
assumes $i > 0$
shows $\{..<i\} \neq \{\}$
 $\langle proof \rangle$

lemma $inj\text{-}on\text{-}inv\text{-}into[simp]$:
assumes $inj\text{-}on\ F\ P$
shows $inv\text{-}into\ P\ F\ ' (F\ ' P) = P$
 $\langle proof \rangle$

lemma $inj\text{-}on\text{-}inv\text{-}into2[simp]$:
 $inj\text{-}on\ (inv\text{-}into\ P\ F)\ (F\ ' P)$
 $\langle proof \rangle$

lemma $refl\text{-}gfp$:
assumes $1: mono\ Retr$ **and** $2: \bigwedge\ theta. refl\ theta \implies refl\ (Retr\ theta)$
shows $refl\ (gfp\ Retr)$
 $\langle proof \rangle$

lemma $sym\text{-}gfp$:
assumes $1: mono\ Retr$ **and** $2: \bigwedge\ theta. sym\ theta \implies sym\ (Retr\ theta)$
shows $sym\ (gfp\ Retr)$
 $\langle proof \rangle$

lemma $trancl\text{-}trans[simp]$:

assumes *trans R*
shows $P \hat{+} \subseteq R \longleftrightarrow P \subseteq R$
 $\langle proof \rangle$

lemma *trans-gfp*:
assumes 1: *mono Retr* **and** 2: $\bigwedge theta. trans\ theta \implies trans\ (Retr\ theta)$
shows $trans\ (gfp\ Retr)$
 $\langle proof \rangle$

lemma *O-subset-trans*:
assumes $r\ O\ r \subseteq r$
shows $trans\ r$
 $\langle proof \rangle$

lemma *trancl-imp-trans*:
assumes $r \hat{+} \subseteq r$
shows $trans\ r$
 $\langle proof \rangle$

lemma *sym-trans-gfp*:
assumes 1: *mono Retr* **and** 2: $\bigwedge theta. sym\ theta \wedge trans\ theta \implies sym\ (Retr\ theta) \wedge trans\ (Retr\ theta)$
shows $sym\ (gfp\ Retr) \wedge trans\ (gfp\ Retr)$
 $\langle proof \rangle$

2.2 Infrastructure for partitions

definition *part where*
 $part\ J\ P \equiv$
 $Union\ P = J \wedge$
 $(\forall J1\ J2. J1 \in P \wedge J2 \in P \wedge J1 \neq J2 \implies J1 \cap J2 = \{\})$

inductive-set *gen*
for $P :: 'a\ set\ set$ **and** $I :: 'a\ set$ **where**
 $incl[simp]: i \in I \implies i \in gen\ P\ I$
 $ext[simp]: \llbracket J \in P; j0 \in J; j0 \in gen\ P\ I; j \in J \rrbracket \implies j \in gen\ P\ I$

definition *partGen where*
 $partGen\ P \equiv \{gen\ P\ I \mid I. I \in P\}$

definition *finer where*
 $finer\ P\ Q \equiv$
 $(\forall J \in Q. J = Union\ \{I \in P. I \subseteq J\}) \wedge$
 $(P \neq \{\} \implies Q \neq \{\})$

definition *partJoin* **where**
partJoin $P Q \equiv \text{partGen } (P \cup Q)$

definition *compat* **where**
compat $I \text{ theta } f \equiv \forall i j. \{i, j\} \subseteq I \wedge i \neq j \longrightarrow (f i, f j) \in \text{theta}$

definition *partCompat* **where**
partCompat $P \text{ theta } f \equiv$
 $\forall I \in P. \text{compat } I \text{ theta } f$

definition *lift* **where**
lift $P F II \equiv \text{Union } \{F I \mid I. I \in P \wedge I \subseteq II\}$

part:

lemma *part-emp[simp]*:
part $J (\text{insert } \{ \} P) = \text{part } J P$
 $\langle \text{proof} \rangle$

lemma *finite-part[simp]*:
assumes *finite* I **and** *part* $I P$
shows *finite* P
 $\langle \text{proof} \rangle$

lemma *part-sum*:
assumes $P: \text{part } \{..<n::\text{nat}\} P$
shows $(\sum i < n. f i) = (\sum p \in P. \sum i \in p. f i)$
 $\langle \text{proof} \rangle$

lemma *part-Un[simp]*:
assumes *part* $I1 P1$ **and** *part* $I2 P2$ **and** $I1 \text{ Int } I2 = \{ \}$
shows *part* $(I1 \text{ Un } I2) (P1 \text{ Un } P2)$
 $\langle \text{proof} \rangle$

lemma *part-Un-singl[simp]*:
assumes *part* $K P$ **and** $\bigwedge I. I \in P \implies I0 \text{ Int } I = \{ \}$
shows *part* $(I0 \text{ Un } K) (\{I0\} \text{ Un } P)$
 $\langle \text{proof} \rangle$

lemma *part-Un-singl2*:
assumes $K01 = I0 \text{ Un } K1$
and *part* $K1 P$ **and** $\bigwedge I. I \in P \implies I0 \text{ Int } I = \{ \}$
shows *part* $K01 (\{I0\} \text{ Un } P)$
 $\langle \text{proof} \rangle$

lemma *part-UN*:
assumes $\bigwedge n. n \in N \implies \text{part } (I n) (P n)$

and $\bigwedge n1\ n2. \{n1,n2\} \subseteq N \wedge n1 \neq n2 \implies I\ n1 \cap I\ n2 = \{\}$
shows $part\ (UN\ n : N. I\ n)\ (UN\ n : N. P\ n)$
 $\langle proof \rangle$

gen:

lemma *incl-gen[simp]*:
 $I \subseteq gen\ P\ I$
 $\langle proof \rangle$

lemma *gen-incl-Un*:
 $gen\ P\ I \subseteq I \cup (Union\ P)$
 $\langle proof \rangle$

lemma *gen-incl*:
assumes $I \in P$
shows $gen\ P\ I \subseteq Union\ P$
 $\langle proof \rangle$

lemma *finite-gen*:
assumes *finite* P **and** $\bigwedge J. J \in P \implies finite\ J$ **and** *finite* I
shows *finite* $(gen\ P\ I)$
 $\langle proof \rangle$

lemma *subset-gen[simp]*:
assumes $J \in P$ **and** $gen\ P\ I \cap J \neq \{\}$
shows $J \subseteq gen\ P\ I$
 $\langle proof \rangle$

lemma *gen-subset-gen[simp]*:
assumes $J \in P$ **and** $gen\ P\ I \cap J \neq \{\}$
shows $gen\ P\ J \subseteq gen\ P\ I$
 $\langle proof \rangle$

lemma *gen-mono[simp]*:
assumes $I \subseteq J$
shows $gen\ P\ I \subseteq gen\ P\ J$
 $\langle proof \rangle$

lemma *gen-idem[simp]*:
 $gen\ P\ (gen\ P\ I) = gen\ P\ I$
 $\langle proof \rangle$

lemma *gen-nchotomy*:
assumes $J \in P$
shows $J \subseteq gen\ P\ I \vee gen\ P\ I \cap J = \{\}$
 $\langle proof \rangle$

lemma *gen-Union*:
assumes $I \in P$

shows $gen\ P\ I = Union\ \{J \in P . J \subseteq gen\ P\ I\}$
 $\langle proof \rangle$

lemma *subset-gen2*:
assumes $*$: $\{I, J\} \subseteq P$ **and** $**$: $gen\ P\ I \cap gen\ P\ J \neq \{\}$
shows $I \subseteq gen\ P\ J$
 $\langle proof \rangle$

lemma *gen-subset-gen2[simp]*:
assumes $\{I, J\} \subseteq P$ **and** $gen\ P\ I \cap gen\ P\ J \neq \{\}$
shows $gen\ P\ I \subseteq gen\ P\ J$
 $\langle proof \rangle$

lemma *gen-eq-gen*:
assumes $\{I, J\} \subseteq P$ **and** $gen\ P\ I \cap gen\ P\ J \neq \{\}$
shows $gen\ P\ I = gen\ P\ J$
 $\langle proof \rangle$

lemma *gen-empty[simp]*:
 $gen\ P\ \{\} = \{\}$
 $\langle proof \rangle$

lemma *gen-empty2[simp]*:
 $gen\ \{\} I = I$
 $\langle proof \rangle$

lemma *emp-gen[simp]*:
assumes $gen\ P\ I = \{\}$
shows $I = \{\}$
 $\langle proof \rangle$

partGen:

lemma *partGen-ex*:
assumes $I \in P$
shows $\exists J \in partGen\ P. I \subseteq J$
 $\langle proof \rangle$

lemma *ex-partGen*:
assumes $J \in partGen\ P$ **and** $j: j \in J$
shows $\exists I \in P. j \in I$
 $\langle proof \rangle$

lemma *Union-partGen*: $\bigcup (partGen\ P) = \bigcup P$
 $\langle proof \rangle$

lemma *Int-partGen*:
assumes $*$: $\{I, J\} \subseteq partGen\ P$ **and** $**$: $I \cap J \neq \{\}$
shows $I = J$
 $\langle proof \rangle$

lemma *part-partGen*:
part (*Union P*) (*partGen P*)
<proof>

lemma *finite-partGen[simp]*:
assumes *finite P*
shows *finite (partGen P)*
<proof>

lemma *emp-partGen[simp]*:
assumes $\{\} \notin P$
shows $\{\} \notin \text{partGen } P$
<proof>

finer:

lemma *finer-partGen*:
finer P (partGen P)
<proof>

lemma *finer-nchotomy*:
assumes *P: part I0 P and Q: part I0 Q and PQ: finer P Q*
and *I: I ∈ P and II: II ∈ Q*
shows $I \subseteq II \vee (I \cap II = \{\})$
<proof>

lemma *finer-ex*:
assumes *P: part I0 P and Q: part I0 Q and PQ: finer P Q*
and *I: I ∈ P*
shows $\exists II. II \in Q \wedge I \subseteq II$
<proof>

partJoin:

lemma *partJoin-commute*:
partJoin P Q = partJoin Q P
<proof>

lemma *Union-partJoin-L*:
 $Union P \subseteq Union (partJoin P Q)$
<proof>

lemma *Union-partJoin-R*:
 $Union Q \subseteq Union (partJoin P Q)$
<proof>

lemma *part-partJoin[simp]*:
assumes *part I P and part I Q*
shows *part I (partJoin P Q)*
<proof>

lemma *finer-partJoin-L[simp]*:
assumes *: *part I P* **and** **: *part I Q*
shows *finer P (partJoin P Q)*
 \langle *proof* \rangle

lemma *finer-partJoin-R[simp]*:
assumes *: *part I P* **and** **: *part I Q*
shows *finer Q (partJoin P Q)*
 \langle *proof* \rangle

lemma *finer-emp[simp]*:
assumes *finer {} Q*
shows $Q \subseteq \{ \{ \} \}$
 \langle *proof* \rangle

compat:

lemma *part-emp-R[simp]*:
part I {} \longleftrightarrow $I = \{ \}$
 \langle *proof* \rangle

lemma *part-emp-L[simp]*:
part {} P \implies $P \subseteq \{ \{ \} \}$
 \langle *proof* \rangle

lemma *finite-partJoin[simp]*:
assumes *finite P* **and** *finite Q*
shows *finite (partJoin P Q)*
 \langle *proof* \rangle

lemma *emp-partJoin[simp]*:
assumes $\{ \} \notin P$ **and** $\{ \} \notin Q$
shows $\{ \} \notin$ *partJoin P Q*
 \langle *proof* \rangle

partCompat:

lemma *partCompat-Un[simp]*:
partCompat (P Un Q) theta f \longleftrightarrow
partCompat P theta f \wedge *partCompat Q theta f*
 \langle *proof* \rangle

lemma *partCompat-gen-aux*:
assumes *theta: sym theta trans theta*
and *fP: partCompat P theta f* **and** $I: I \in P$
and $i: i \in I$ **and** $j: j \in$ *gen P I* **and** $ij: i \neq j$
shows $(f\ i, f\ j) \in$ *theta*
 \langle *proof* \rangle

lemma *partCompat-gen*:

assumes *theta: sym theta trans theta*
and *fP: partCompat P theta f* **and** *I: I ∈ P*
shows *compat (gen P I) theta f*
<proof>

lemma *partCompat-partGen:*
assumes *sym theta* **and** *trans theta*
and *partCompat P theta f*
shows *partCompat (partGen P) theta f*
<proof>

lemma *partCompat-partJoin[simp]:*
assumes *sym theta* **and** *trans theta*
and *partCompat P theta f* **and** *partCompat Q theta f*
shows *partCompat (partJoin P Q) theta f*
<proof>

lift:

lemma *inj-on-lift:*
assumes *P: part I0 P* **and** *Q: part I0 Q* **and** *PQ: finer P Q*
and *F: inj-on F P* **and** *FP: part J0 (F ◁ P)* **and** *emp: {} ∉ F ◁ P*
shows *inj-on (lift P F) Q*
<proof>

lemma *part-lift:*
assumes *P: part I0 P* **and** *Q: part I0 Q* **and** *PQ: finer P Q*
and *F: inj-on F P* **and** *FP: part J0 (F ◁ P)* **and** *emp: {} ∉ P {} ∉ F ◁ P*
shows *part J0 (lift P F ◁ Q)*
<proof>

lemma *finer-lift:*
assumes *finer P Q*
shows *finer (F ◁ P) (lift P F ◁ Q)*
<proof>

2.3 Basic setting for bisimilarity

locale *PL-Indis =*
PL aval tval cval
for *aval :: 'atom ⇒ 'state ⇒ 'state* **and**
tval :: 'test ⇒ 'state ⇒ bool **and**
cval :: 'choice ⇒ 'state ⇒ real +
fixes
indis :: 'state rel
assumes
equiv-indis: equiv UNIV indis

context *PL-Indis*
begin

no-notation *eqpoll* (**infixl** $\langle \approx \rangle$ 50)

abbreviation *indisAbbrev* (**infix** $\langle \approx \rangle$ 50)
where $s1 \approx s2 \equiv (s1, s2) \in indis$

lemma *refl-indis*: *refl indis*
and *trans-indis*: *trans indis*
and *sym-indis*: *sym indis*
\langle proof \rangle

lemma *indis-refl*[*intro*]: $s \approx s$
\langle proof \rangle

lemma *indis-trans*[*trans*]: $\llbracket s \approx s'; s' \approx s'' \rrbracket \implies s \approx s''$
\langle proof \rangle

lemma *indis-sym*[*sym*]: $s \approx s' \implies s' \approx s$
\langle proof \rangle

2.4 Discreetness

coinductive *discr* **where**

intro:
 $(\bigwedge s i. i < brn\ c \longrightarrow s \approx eff\ c\ s\ i \wedge discr\ (cont\ c\ s\ i))$
 $\implies discr\ c$

definition *discrL* **where**

$discrL\ cl \equiv \forall c \in set\ cl. discr\ c$

lemma *discrL-intro*[*intro*]:
assumes $\bigwedge c. c \in set\ cl \implies discr\ c$
shows *discrL cl*
\langle proof \rangle

lemma *discrL-discr*[*simp*]:
assumes *discrL cl* **and** $c \in set\ cl$
shows *discr c*
\langle proof \rangle

lemma *discrL-update*[*simp*]:
assumes *cl: discrL cl* **and** $c': discr\ c'$
shows *discrL (cl[n := c'])*
\langle proof \rangle

Coinduction for discreetness:

lemma *discr-coind*[*consumes 1, case-names Hyp, induct pred: discr*]:

assumes *: $\text{phi } c$ **and**
 **: $\bigwedge c s i. [\text{phi } c ; i < \text{brn } c]$
 $\implies s \approx \text{eff } c s i \wedge (\text{phi } (\text{cont } c s i) \vee \text{discr } (\text{cont } c s i))$
shows $\text{discr } c$
 $\langle \text{proof} \rangle$

lemma discr-raw-coind [$\text{consumes } 1, \text{case-names Hyp}$]:
assumes *: $\text{phi } c$ **and**
 **: $\bigwedge c s i. [i < \text{brn } c; \text{phi } c] \implies s \approx \text{eff } c s i \wedge \text{phi } (\text{cont } c s i)$
shows $\text{discr } c$
 $\langle \text{proof} \rangle$

Discreetness versus transition:

lemma discr-cont [simp]:
assumes *: $\text{discr } c$ **and** **: $i < \text{brn } c$
shows $\text{discr } (\text{cont } c s i)$
 $\langle \text{proof} \rangle$

lemma discr-eff-indis [simp]:
assumes *: $\text{discr } c$ **and** **: $i < \text{brn } c$
shows $s \approx \text{eff } c s i$
 $\langle \text{proof} \rangle$

2.5 Self-isomorphism

coinductive siso **where**

intro:

$[[\bigwedge s i. i < \text{brn } c \implies \text{siso } (\text{cont } c s i);$
 $\bigwedge s t i.$
 $i < \text{brn } c \wedge s \approx t \implies$
 $\text{eff } c s i \approx \text{eff } c t i \wedge \text{wt } c s i = \text{wt } c t i \wedge \text{cont } c s i = \text{cont } c t i]]$
 $\implies \text{siso } c$

definition sisoL **where**

$\text{sisoL } cl \equiv \forall c \in \text{set } cl. \text{siso } c$

lemma sisoL-intro [intro]:
assumes $\bigwedge c. c \in \text{set } cl \implies \text{siso } c$
shows $\text{sisoL } cl$
 $\langle \text{proof} \rangle$

lemma sisoL-siso [simp]:
assumes $\text{sisoL } cl$ **and** $c \in \text{set } cl$
shows $\text{siso } c$
 $\langle \text{proof} \rangle$

lemma sisoL-update [simp]:
assumes $cl: \text{sisoL } cl$ **and** $c': \text{siso } c'$
shows $\text{sisoL } (cl[n := c'])$

$\langle proof \rangle$

Coinduction for self-isomorphism:

lemma *siso-coind*[*consumes 1, case-names Obs Cont, induct pred: siso*]:

assumes *: *phi c* **and**

** : $\bigwedge c s t i. \llbracket i < brn c; phi c; s \approx t \rrbracket \implies$

$eff c s i \approx eff c t i \wedge wt c s i = wt c t i \wedge cont c s i = cont c t i$ **and**

*** : $\bigwedge c s i. \llbracket i < brn c; phi c \rrbracket \implies phi (cont c s i) \vee siso (cont c s i)$

shows *siso c*

$\langle proof \rangle$

lemma *siso-raw-coind*[*consumes 1, case-names Obs Cont*]:

assumes *: *phi c* **and**

*** : $\bigwedge c s t i. \llbracket i < brn c; phi c; s \approx t \rrbracket \implies$

$eff c s i \approx eff c t i \wedge wt c s i = wt c t i \wedge cont c s i = cont c t i$ **and**

** : $\bigwedge c s i. \llbracket i < brn c; phi c \rrbracket \implies phi (cont c s i)$

shows *siso c*

$\langle proof \rangle$

Self-Isomorphism versus transition:

lemma *siso-cont*[*simp*]:

assumes *: *siso c* **and** **: $i < brn c$

shows *siso (cont c s i)*

$\langle proof \rangle$

lemma *siso-cont-indis*[*simp*]:

assumes *: *siso c* **and** **: $s \approx t i < brn c$

shows $eff c s i \approx eff c t i \wedge wt c s i = wt c t i \wedge cont c s i = cont c t i$

$\langle proof \rangle$

2.6 Notions of bisimilarity

Matchers

definition *mC-C-part* **where**

mC-C-part $c d P F \equiv$

$\{\} \notin P \wedge \{\} \notin F \text{ ' } P \wedge$

part $\{.. < brn c\} P \wedge \text{part } \{.. < brn d\} (F \text{ ' } P)$

definition *mC-C-wt* **where**

mC-C-wt $c d s t P F \equiv \forall I \in P. \text{sum } (wt c s) I = \text{sum } (wt d t) (F I)$

definition *mC-C-eff-cont* **where**

mC-C-eff-cont $theta c d s t P F \equiv$

$\forall I i j.$

$I \in P \wedge i \in I \wedge j \in F I \longrightarrow$

$eff c s i \approx eff d t j \wedge (cont c s i, cont d t j) \in theta$

definition *mC-C* **where**

$mC-C$ $theta$ c d s t P F \equiv
 $mC-C-part$ c d P F $\wedge inj-on$ F P $\wedge mC-C-wt$ c d s t P F $\wedge mC-C-eff-cont$ $theta$
 c d s t P F

definition $matchC-C$ **where**

$matchC-C$ $theta$ c d $\equiv \forall s t. s \approx t \longrightarrow (\exists P F. mC-C$ $theta$ c d s t P F)

definition $mC-ZOC-part$ **where**

$mC-ZOC-part$ c d s t $I0$ P F \equiv
 $\{\}$ $\notin P - \{I0\} \wedge \{\}$ $\notin F '(P - \{I0\}) \wedge I0 \in P \wedge$
 $part$ $\{..< brn c\}$ $P \wedge part$ $\{..< brn d\}$ $(F ' P)$

definition $mC-ZOC-wt$ **where**

$mC-ZOC-wt$ c d s t $I0$ P F \equiv
 sum $(wt$ c $s)$ $I0 < 1 \wedge sum$ $(wt$ d $t)$ $(F$ $I0) < 1 \longrightarrow$
 $(\forall I \in P - \{I0\}.$
 sum $(wt$ c $s)$ $I / (1 - sum$ $(wt$ c $s)$ $I0) =$
 sum $(wt$ d $t)$ $(F$ $I) / (1 - sum$ $(wt$ d $t)$ $(F$ $I0)))$

definition $mC-ZOC-eff-cont0$ **where**

$mC-ZOC-eff-cont0$ $theta$ c d s t $I0$ F \equiv
 $(\forall i \in I0. s \approx eff$ c s $i \wedge (cont$ c s $i, d) \in theta) \wedge$
 $(\forall j \in F$ $I0. t \approx eff$ d t $j \wedge (c, cont$ d t $j) \in theta)$

definition $mC-ZOC-eff-cont$ **where**

$mC-ZOC-eff-cont$ $theta$ c d s t $I0$ P F \equiv
 $\forall I$ i $j.$
 $I \in P - \{I0\} \wedge i \in I \wedge j \in F$ $I \longrightarrow$
 eff c s $i \approx eff$ d t $j \wedge$
 $(cont$ c s $i, cont$ d t $j) \in theta$

definition $mC-ZOC$ **where**

$mC-ZOC$ $theta$ c d s t $I0$ P F \equiv
 $mC-ZOC-part$ c d s t $I0$ P F \wedge
 $inj-on$ F P \wedge
 $mC-ZOC-wt$ c d s t $I0$ P F \wedge
 $mC-ZOC-eff-cont0$ $theta$ c d s t $I0$ F \wedge
 $mC-ZOC-eff-cont$ $theta$ c d s t $I0$ P F

definition $matchC-LC$ **where**

$matchC-LC$ $theta$ c d \equiv
 $\forall s t. s \approx t \longrightarrow (\exists I0$ P $F. mC-ZOC$ $theta$ c d s t $I0$ P $F)$

lemmas $m-defs = mC-C-def$ $mC-ZOC-def$

lemmas $m-defsAll =$

$mC-C-def$ $mC-C-part-def$ $mC-C-wt-def$ $mC-C-eff-cont-def$

mC-ZOC-def mC-ZOC-part-def mC-ZOC-wt-def mC-ZOC-eff-cont0-def mC-ZOC-eff-cont-def

lemmas *match-defs* =
matchC-C-def matchC-LC-def

lemma *mC-C-mono*:
assumes *mC-C theta c d s t P F* **and** $\theta \subseteq \theta'$
shows *mC-C theta' c d s t P F*
<proof>

lemma *matchC-C-mono*:
assumes *matchC-C theta c d* **and** $\theta \subseteq \theta'$
shows *matchC-C theta' c d*
<proof>

lemma *mC-ZOC-mono*:
assumes *mC-ZOC theta c d s t I0 P F* **and** $\theta \subseteq \theta'$
shows *mC-ZOC theta' c d s t I0 P F*
<proof>

lemma *matchC-LC-mono*:
assumes *matchC-LC theta c d* **and** $\theta \subseteq \theta'$
shows *matchC-LC theta' c d*
<proof>

lemma *Int-not-in-eq-emp*:
 $P \cap \{I. I \notin P\} = \{\}$
<proof>

lemma *mC-C-mC-ZOC*:
assumes *mC-C theta c d s t P F*
shows *mC-ZOC theta c d s t {} (P Un { {} }) (%I. if I ∈ P then F I else {})*
(**is** *mC-ZOC theta c d s t ?I0 ?Q ?G*)
<proof>

lemma *matchC-C-matchC-LC*:
assumes *matchC-C theta c d*
shows *matchC-LC theta c d*
<proof>

Retracts:

definition *Sretr* **where**
Sretr theta \equiv
 $\{(c, d). \text{matchC-C } \theta \ c \ d\}$

definition *ZOretr* **where**
ZOretr theta \equiv
 $\{(c, d). \text{matchC-LC } \theta \ c \ d\}$

lemmas *Retr-defs* =
Sretr-def
ZOretr-def

lemma *mono-Retr*:
mono Sretr
mono ZOretr
(*proof*)

lemma *Retr-incl*:
Sretr theta \subseteq *ZOretr theta*
(*proof*)

The associated bisimilarity relations:

definition *Sbis* **where** *Sbis* \equiv *gfp Sretr*

definition *ZObis* **where** *ZObis* \equiv *gfp ZOretr*

abbreviation *Sbis-abbrev* (**infix** $\langle \approx s \rangle$ 55) **where** $c \approx s d \equiv (c, d) : Sbis$

abbreviation *ZObis-abbrev* (**infix** $\langle \approx 01 \rangle$ 55) **where** $c \approx 01 d \equiv (c, d) : ZObis$

lemmas *bis-defs* = *Sbis-def ZObis-def*

lemma *bis-incl*:
Sbis \leq *ZObis*
(*proof*)

lemma *bis-imp[simp]*:
 $\bigwedge c1 c2. c1 \approx s c2 \implies c1 \approx 01 c2$
(*proof*)

lemma *Sbis-prefix*:
Sbis \leq *Sretr Sbis*
(*proof*)

lemma *Sbis-fix*:
Sretr Sbis = *Sbis*
(*proof*)

lemma *Sbis-mC-C*:
assumes $c \approx s d$ **and** $s \approx t$
shows $\exists P F. mC-C Sbis c d s t P F$
(*proof*)

lemma *Sbis-coind*:
assumes $theta \leq Sretr$ (*theta Un Sbis*)

shows $\theta \leq Sbis$
<proof>

lemma *Sbis-raw-coind*:
assumes $\theta \leq Sretr \theta$
shows $\theta \leq Sbis$
<proof>

lemma *mC-C-sym*:
assumes $mC-C \theta c d s t P F$
shows $mC-C (\theta^{-1}) d c t s (F \text{ ' } P) (inv\text{-}into P F)$
<proof>

lemma *matchC-C-sym*:
assumes $matchC-C \theta c d$
shows $matchC-C (\theta^{-1}) d c$
<proof>

lemma *Sretr-sym*:
assumes $sym \theta$
shows $sym (Sretr \theta)$
<proof>

lemma *sym-Sbis*: $sym Sbis$
<proof>

lemma *Sbis-sym*: $c \approx s d \implies d \approx s c$
<proof>

lemma *mC-C-trans*:
assumes $*$: $mC-C \theta_1 c d s t P F$ **and** $*$: $mC-C \theta_2 d e t u (F \text{ ' } P) G$
shows $mC-C (\theta_1 \circ \theta_2) c e s u P (G \circ F)$
<proof>

lemma *mC-C-finer*:
assumes $*$: $mC-C \theta c d s t P F$
and θ : $trans \theta$
and Q : $finer P Q \text{ finite } Q \} \notin Q \text{ part } \{..<brn c\} Q$
and c : $partCompat Q indis (eff c s) partCompat Q \theta (cont c s)$
shows $mC-C \theta c d s t Q (lift P F)$
<proof>

lemma *mC-C-partCompat-eff*:
assumes $*$: $mC-C \theta c d s t P F$

shows *partCompat P indis (eff c s)*
{proof}

lemma *mC-C-partCompat-cont*:
assumes *: *mC-C theta c d s t P F*
and *theta: sym theta trans theta*
shows *partCompat P theta (cont c s)*
{proof}

lemma *matchC-C-sym-trans*:
assumes *: *matchC-C theta c1 c* **and** **: *matchC-C theta c c2*
and *theta: sym theta trans theta*
shows *matchC-C theta c1 c2*
{proof}

lemma *Sretr-sym-trans*:
assumes *sym theta \wedge trans theta*
shows *trans (Sretr theta)*
{proof}

lemma *trans-Sbis: trans Sbis*
{proof}

lemma *Sbis-trans: $\llbracket c \approx_s d; d \approx_s e \rrbracket \implies c \approx_s e$*
{proof}

lemma *ZObis-prefix*:
ZObis \leq ZOretr ZObis
{proof}

lemma *ZObis-fix*:
ZOretr ZObis = ZObis
{proof}

lemma *ZObis-mC-ZOC*:
assumes *c \approx_{01} d* **and** *s \approx t*
shows $\exists I0 P F. mC-ZOC ZObis c d s t I0 P F$
{proof}

lemma *ZObis-coind*:
assumes *theta \leq ZOretr (theta Un ZObis)*
shows *theta \leq ZObis*
{proof}

lemma *ZObis-raw-coind*:
assumes *theta \leq ZOretr theta*

shows $\theta \leq ZObis$
<proof>

lemma *mC-ZOC-sym*:
assumes θ : *sym* θ **and** $*$: *mC-ZOC* θ c d s t $I0$ P F
shows *mC-ZOC* θ d c t s $(F$ $I0)$ $(F$ $'$ $P)$ $(inv$ -*into* P $F)$
<proof>

lemma *matchC-LC-sym*:
assumes $*$: *sym* θ **and** *matchC-LC* θ c d
shows *matchC-LC* θ d c
<proof>

lemma *ZOretr-sym*:
assumes *sym* θ
shows *sym* $(ZOretr$ $\theta)$
<proof>

lemma *sym-ZObis*: *sym* $ZObis$
<proof>

lemma *ZObis-sym*: $c \approx_{01} d \implies d \approx_{01} c$
<proof>

2.7 List versions of the bisimilarities

definition *SbisL* where
 $SbisL$ cl $dl \equiv$
 $length$ $cl = length$ $dl \wedge (\forall n < length$ $cl. cl ! n \approx_s dl ! n)$

lemma *SbisL-intro*[*intro*]:
assumes $length$ $cl = length$ dl **and**
 $\bigwedge n. \llbracket n < length$ $cl; n < length$ $dl \rrbracket \implies cl ! n \approx_s dl ! n$
shows *SbisL* cl dl
<proof>

lemma *SbisL-length*[*simp*]:
assumes *SbisL* cl dl
shows $length$ $cl = length$ dl
<proof>

lemma *SbisL-Sbis*[*simp*]:
assumes *SbisL* cl dl **and** $n < length$ $cl \vee n < length$ dl
shows $cl ! n \approx_s dl ! n$
<proof>

lemma *SbisL-update*[*simp*]:

assumes $cldl$: $SbisL\ cl\ dl$ **and** $c'd'$: $c' \approx_s d'$
shows $SbisL\ (cl\ [n := c'])\ (dl\ [n := d'])$ (**is** $SbisL\ ?cl'\ ?dl'$)
 $\langle proof \rangle$

lemma $SbisL\text{-update-L}[simp]$:
assumes $SbisL\ cl\ dl$ **and** $c' \approx_s dl!n$
shows $SbisL\ (cl[n := c'])\ dl$
 $\langle proof \rangle$

lemma $SbisL\text{-update-R}[simp]$:
assumes $SbisL\ cl\ dl$ **and** $cl!n \approx_s d'$
shows $SbisL\ cl\ (dl[n := d'])$
 $\langle proof \rangle$

definition $ZObisL$ **where**
 $ZObisL\ cl\ dl \equiv$
 $length\ cl = length\ dl \wedge (\forall\ n < length\ cl.\ cl!\ n \approx_{01}\ dl!\ n)$

lemma $ZObisL\text{-intro}[intro]$:
assumes $length\ cl = length\ dl$ **and**
 $\bigwedge\ n.\ [n < length\ cl; n < length\ dl] \implies cl!\ n \approx_{01}\ dl!\ n$
shows $ZObisL\ cl\ dl$
 $\langle proof \rangle$

lemma $ZObisL\text{-length}[simp]$:
assumes $ZObisL\ cl\ dl$
shows $length\ cl = length\ dl$
 $\langle proof \rangle$

lemma $ZObisL\text{-ZObis}[simp]$:
assumes $ZObisL\ cl\ dl$ **and** $n < length\ cl \vee n < length\ dl$
shows $cl!\ n \approx_{01}\ dl!\ n$
 $\langle proof \rangle$

lemma $ZObisL\text{-update}[simp]$:
assumes $cldl$: $ZObisL\ cl\ dl$ **and** $c'd'$: $c' \approx_{01}\ d'$
shows $ZObisL\ (cl\ [n := c'])\ (dl\ [n := d'])$ (**is** $ZObisL\ ?cl'\ ?dl'$)
 $\langle proof \rangle$

lemma $ZObisL\text{-update-L}[simp]$:
assumes $ZObisL\ cl\ dl$ **and** $c' \approx_{01}\ dl!n$
shows $ZObisL\ (cl[n := c'])\ dl$
 $\langle proof \rangle$

lemma $ZObisL\text{-update-R}[simp]$:
assumes $ZObisL\ cl\ dl$ **and** $cl!n \approx_{01}\ d'$
shows $ZObisL\ cl\ (dl[n := d'])$

<proof>

2.8 Discreetness for configurations

coinductive *discrCf* **where**

intro:

$(\bigwedge i. i < \text{brn} (\text{fst } cf) \longrightarrow$
 $\text{snd } cf \approx \text{snd} (\text{cont-eff } cf \ i) \wedge \text{discrCf} (\text{cont-eff } cf \ i)$
 $\implies \text{discrCf } cf$

Coinduction for discrness:

lemma *discrCf-coind*[*consumes 1, case-names Hyp, induct pred: discr*]:

assumes *: *phi cf* **and**

** : $\bigwedge cf \ i.$

$\llbracket i < \text{brn} (\text{fst } cf); \text{phi } cf \rrbracket \implies$

$\text{snd } cf \approx \text{snd} (\text{cont-eff } cf \ i) \wedge (\text{phi} (\text{cont-eff } cf \ i) \vee \text{discrCf} (\text{cont-eff } cf \ i))$

shows *discrCf cf*

<proof>

lemma *discrCf-raw-coind*[*consumes 1, case-names Hyp*]:

assumes *: *phi cf* **and**

** : $\bigwedge cf \ i. \llbracket i < \text{brn} (\text{fst } cf); \text{phi } cf \rrbracket \implies$

$\text{snd } cf \approx \text{snd} (\text{cont-eff } cf \ i) \wedge \text{phi} (\text{cont-eff } cf \ i)$

shows *discrCf cf*

<proof>

Discreetness versus transition:

lemma *discrCf-cont*[*simp*]:

assumes *: *discrCf cf* **and** **: $i < \text{brn} (\text{fst } cf)$

shows *discrCf (cont-eff cf i)*

<proof>

lemma *discrCf-eff-indis*[*simp*]:

assumes *: *discrCf cf* **and** **: $i < \text{brn} (\text{fst } cf)$

shows $\text{snd } cf \approx \text{snd} (\text{cont-eff } cf \ i)$

<proof>

lemma *discr-discrCf*:

assumes *discr c*

shows *discrCf (c, s)*

<proof>

lemma *ZObis-pres-discrCfL*:

assumes $\text{fst } cf \approx_{01} \text{fst } df$ **and** $\text{snd } cf \approx \text{snd } df$ **and** *discrCf df*

shows *discrCf cf*

<proof>

corollary *ZObis-pres-discrCfR*:

assumes *discrCf cf* **and** $\text{fst } cf \approx_{01} \text{fst } df$ **and** $\text{snd } cf \approx \text{snd } df$

shows $\text{discrCf } df$
 $\langle \text{proof} \rangle$

end

end

3 Trace-Based Noninterference

theory *Trace-Based*
 imports *Resumption-Based*
begin

3.1 Preliminaries

lemma *dist-sum*:

fixes $f :: 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow \text{real}$
 assumes $\bigwedge i. i \in I \Longrightarrow \text{dist } (f \ i) \ (g \ i) \leq e \ i$
 shows $\text{dist } (\sum i \in I. f \ i) \ (\sum i \in I. g \ i) \leq (\sum i \in I. e \ i)$
 $\langle \text{proof} \rangle$

lemma *dist-mult[simp]*: $\text{dist } (x * y) \ (x * z) = |x| * \text{dist } y \ z$
 $\langle \text{proof} \rangle$

lemma *dist-divide[simp]*: $\text{dist } (y / r) \ (z / r) = \text{dist } y \ z / |r|$
 $\langle \text{proof} \rangle$

lemma *dist-weighted-sum*:

fixes $f :: 'a \Rightarrow \text{real}$ **and** $g :: 'b \Rightarrow \text{real}$
 assumes $\text{eps}: \bigwedge i \ j. i \in I \Longrightarrow j \in J \Longrightarrow w \ i \neq 0 \Longrightarrow v \ j \neq 0 \Longrightarrow \text{dist } (f \ i) \ (g \ j) \leq d \ i + e \ j$
 and $\text{pos}: \bigwedge i. i \in I \Longrightarrow 0 \leq w \ i \ \bigwedge j. j \in J \Longrightarrow 0 \leq v \ j$
 and $\text{sum}: (\sum i \in I. w \ i) = 1 \ (\sum j \in J. v \ j) = 1$
 shows $\text{dist } (\sum i \in I. w \ i * f \ i) \ (\sum j \in J. v \ j * g \ j) \leq (\sum i \in I. w \ i * d \ i) + (\sum j \in J. v \ j * e \ j)$
 $\langle \text{proof} \rangle$

lemma *field-abs-le-zero-epsilon*:

fixes $x :: 'a::\{\text{linordered-field}\}$
 assumes $\text{epsilon}: \bigwedge e. 0 < e \Longrightarrow |x| \leq e$
 shows $|x| = 0$
 $\langle \text{proof} \rangle$

lemma *nat-nat-induct[case-names less]*:

fixes $P :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{bool}$
 assumes $\text{less}: \bigwedge n \ m. (\bigwedge j \ k. j + k < n + m \Longrightarrow P \ j \ k) \Longrightarrow P \ n \ m$

shows $P \ n \ m$
 $\langle proof \rangle$

lemma *part-insert*:
assumes *part* $A \ P$ **assumes** $X \cap A = \{\}$
shows *part* $(A \cup X)$ (*insert* $X \ P$)
 $\langle proof \rangle$

lemma *part-insert-subset*:
assumes $X: \textit{part} (A - X) \ P \ X \subseteq A$
shows *part* A (*insert* $X \ P$)
 $\langle proof \rangle$

lemma *part-is-subset*:
 $\textit{part} \ S \ P \implies p \in P \implies p \subseteq S$
 $\langle proof \rangle$

lemma *dist-nonneg-bounded*:
fixes $l \ u \ x \ y :: \textit{real}$
assumes $l \leq x \ x \leq u \ l \leq y \ y \leq u$
shows $\textit{dist} \ x \ y \leq u - l$
 $\langle proof \rangle$

lemma *integrable-count-space-finite-support*:
fixes $f :: 'a \Rightarrow 'b :: \{\textit{banach}, \textit{second-countable-topology}\}$
shows $\textit{finite} \ \{x \in X. \ f \ x \neq 0\} \implies \textit{integrable} \ (\textit{count-space} \ X) \ f$
 $\langle proof \rangle$

lemma *lebesgue-integral-point-measure*:
fixes $g :: - \Rightarrow \textit{real}$
assumes $f: \textit{finite} \ \{a \in A. \ 0 < f \ a \wedge g \ a \neq 0\}$
shows $\textit{integral}^L \ (\textit{point-measure} \ A \ f) \ g = (\sum \ a | a \in A \wedge 0 < f \ a \wedge g \ a \neq 0. \ f \ a * g \ a)$
 $\langle proof \rangle$

lemma (**in** *finite-measure*) *finite-measure-dist*:
assumes $AE: AE \ x \ \textit{in} \ M. \ x \notin C \longrightarrow (x \in A \longleftrightarrow x \in B)$
assumes $\textit{sets}: A \in \textit{sets} \ M \ B \in \textit{sets} \ M \ C \in \textit{sets} \ M$
shows $\textit{dist} \ (\textit{measure} \ M \ A) \ (\textit{measure} \ M \ B) \leq \textit{measure} \ M \ C$
 $\langle proof \rangle$

lemma (**in** *prob-space*) *prob-dist*:
assumes $AE: AE \ x \ \textit{in} \ M. \ \neg \ C \ x \longrightarrow (A \ x \longleftrightarrow B \ x)$
assumes $\textit{sets}: \textit{Measurable.pred} \ M \ A \ \textit{Measurable.pred} \ M \ B \ \textit{Measurable.pred} \ M \ C$
shows $\textit{dist} \ \mathcal{P}(x \ \textit{in} \ M. \ A \ x) \ \mathcal{P}(x \ \textit{in} \ M. \ B \ x) \leq \mathcal{P}(x \ \textit{in} \ M. \ C \ x)$
 $\langle proof \rangle$

lemma *Least-eq-0-iff*: $(\exists i :: \textit{nat}. \ P \ i) \implies (\textit{LEAST} \ i. \ P \ i) = 0 \longleftrightarrow P \ 0$

<proof>

lemma *case-nat-comp-Suc*[simp]: $\text{case-nat } x f \circ \text{Suc} = f$
<proof>

lemma *sum-eq-0-iff*:

fixes $f :: - \Rightarrow 'a :: \{\text{comm-monoid-add, ordered-ab-group-add}\}$
shows $\text{finite } A \Longrightarrow (\bigwedge i. i \in A \Longrightarrow 0 \leq f i) \Longrightarrow (\sum_{i \in A}. f i) = 0 \longleftrightarrow (\forall i \in A. f i = 0)$
<proof>

lemma *sum-less-0-iff*:

fixes $f :: - \Rightarrow 'a :: \{\text{comm-monoid-add, ordered-ab-group-add}\}$
shows $\text{finite } A \Longrightarrow (\bigwedge i. i \in A \Longrightarrow 0 \leq f i) \Longrightarrow 0 < (\sum_{i \in A}. f i) \longleftrightarrow (\exists i \in A. 0 < f i)$
<proof>

context *PL-Indis*
begin

declare *emp-gen*[simp del]

interpretation *pmf-as-function* *<proof>*

lift-definition *wt-pmf* :: $('test, 'atom, 'choice) \text{ cmd} \times 'state \Rightarrow \text{nat pmf is}$
 $\lambda(c, s) i. \text{if proper } c \text{ then if } i < \text{brn } c \text{ then wt } c \ s \ i \text{ else } 0 \text{ else if } i = 0 \text{ then } 1 \text{ else } 0$
<proof>

definition *trans* :: $('test, 'atom, 'choice) \text{ cmd} \times 'state \Rightarrow (('test, 'atom, 'choice) \text{ cmd} \times 'state) \text{ pmf where}$
 $\text{trans } cf = \text{map-pmf } (\lambda i. \text{if proper } (fst \ cf) \text{ then cont-eff } cf \ i \text{ else } cf) \ (wt\text{-pmf } cf)$

sublocale *T?*: *MC-syntax trans* *<proof>*

abbreviation $G \ cf \equiv \text{set-pmf } (trans \ cf)$

lemma *set-pmf-map*: $\text{set-pmf } (\text{map-pmf } f \ M) = f \ ' \ \text{set-pmf } M$
<proof>

lemma *set-pmf-wt*:

$\text{set-pmf } (wt\text{-pmf } cf) = (\text{if proper } (fst \ cf) \text{ then } \{i. i < \text{brn } (fst \ cf) \wedge 0 < wt \ (fst \ cf) \ (snd \ cf) \ i\} \text{ else } \{0\})$
<proof>

lemma *G-eq*:

$G \ cf = (\text{if proper } (fst \ cf) \text{ then } \{\text{cont-eff } cf \ i \mid i. i < \text{brn } (fst \ cf) \wedge 0 < wt \ (fst \ cf) \ (snd \ cf) \ i\} \text{ else } \{cf\})$
<proof>

lemma *discrCf-G*: $\text{discrCf } cf \implies cf' \in G \text{ } cf \implies \text{discrCf } cf'$

<proof>

lemma *measurable-discrCf[measurable]*: $\text{Measurable.pred (count-space UNIV) discrCf}$

<proof>

lemma *measurable-indis[measurable]*: $\text{Measurable.pred (count-space UNIV) } (\lambda x. \text{snd } x \approx c)$

<proof>

lemma *integral-trans*:

proper (fst cf) \implies

$(\int x. f x \partial \text{trans } cf) = (\sum i < \text{brn } (\text{fst } cf). \text{wt } (\text{fst } cf) (\text{snd } cf) i * f (\text{cont-eff } cf i))$

<proof>

3.2 Quasi strong termination traces

abbreviation *qsend* \equiv *sfirst (holds discrCf)*

lemma *qsend-eq-0-iff*: $\text{qsend } cfT = 0 \iff \text{discrCf (shd } cfT)$

<proof>

lemma *qsend-eq-0[simp]*: $\text{discrCf } cf \implies \text{qsend } (cf \#\#\ \omega) = 0$

<proof>

lemma *alw-discrCf*: $\text{enabled } cf \ \omega \implies \text{discrCf } cf \implies \text{alw (holds discrCf) } \omega$

<proof>

lemma *alw-discrCf-indis'*:

$\text{enabled } cf \ \omega \implies \text{discrCf } cf \implies \text{snd } cf \approx t \implies \text{alw (holds } (\lambda cf'. \text{snd } cf' \approx t)) \ \omega$

<proof>

lemma *alw-discrCf-indis*:

$\text{enabled } cf \ \omega \implies \text{discrCf } cf \implies \text{alw (holds } (\lambda cf'. \text{snd } cf' \approx \text{snd } cf)) (cf \#\#\ \omega)$

<proof>

lemma *enabled-sdrop*: $\text{enabled } cf \ \omega \implies \text{enabled } ((cf \#\#\ \omega) \#\ n) (\text{sdrop } n \ \omega)$

<proof>

lemma *sfirst-eq-eSuc*: $\text{sfirst } P \ \omega = \text{eSuc } n \iff (\neg P \ \omega) \wedge \text{sfirst } P (\text{stl } \omega) = n$

<proof>

lemma *qsend-snth*: $\text{qsend } \omega = \text{enat } n \implies \text{discrCf } (\omega \#\ n)$

<proof>

lemma *indis-iff*: $a \approx d \implies b \approx d \implies a \approx c \iff b \approx c$

$\langle \text{proof} \rangle$

lemma *enabled-qsend-indis:*

assumes *enabled cf ω qsend (cf ## ω) $\leq n$ qsend (cf ## ω) $\leq m$*

shows *snd ((cf ## ω) !! n) $\approx t \longleftrightarrow$ snd ((cf ## ω) !! m) $\approx t$*

$\langle \text{proof} \rangle$

lemma *enabled-imp-UNTIL-alw-discrCf:*

enabled (shd ω) (stl ω) \implies (not (holds discrCf) until (alw (holds discrCf))) ω

$\langle \text{proof} \rangle$

lemma *less-qsend-iff-not-discrCf:*

enabled cf bT $\implies n < \text{qsend (cf ## bT)} \longleftrightarrow \neg \text{discrCf ((cf ## bT) !! n)$

$\langle \text{proof} \rangle$

3.3 Terminating configurations

definition *qstermCf cf $\longleftrightarrow (\forall cfT. \text{enabled cf cfT} \longrightarrow \text{qsend (cf ## cfT)} < \infty)$*

lemma *qstermCf-E:*

qstermCf cf $\implies cf' \in G \text{ cf} \implies \text{qstermCf cf}'$

$\langle \text{proof} \rangle$

abbreviation *eff-at cf bT n $\equiv \text{snd ((cf ## bT) !! n)$*

abbreviation *cont-at cf bT n $\equiv \text{fst ((cf ## bT) !! n)$*

definition *amSec c $\longleftrightarrow (\forall s1 s2 n t. s1 \approx s2 \longrightarrow$*

$\mathcal{P}(bT \text{ in } T.T (c, s1). \text{eff-at (c, s1) bT n} \approx t) =$

$\mathcal{P}(bT \text{ in } T.T (c, s2). \text{eff-at (c, s2) bT n} \approx t)$

definition *eSec c $\longleftrightarrow (\forall s1 s2 t. s1 \approx s2 \longrightarrow$*

$\mathcal{P}(bT \text{ in } T.T (c, s1). \exists n. \text{qsend ((c, s1) ## bT)} = n \wedge \text{eff-at (c, s1) bT n} \approx t) =$

$\mathcal{P}(bT \text{ in } T.T (c, s2). \exists n. \text{qsend ((c, s2) ## bT)} = n \wedge \text{eff-at (c, s2) bT n} \approx t)$

definition *aeT c $\longleftrightarrow (\forall s. AE bT \text{ in } T.T (c,s). \text{qsend ((c,s) ## bT)} < \infty)$*

lemma *dist-Ps-upper-bound:*

fixes *cf1 cf2 :: ('test, 'atom, 'choice) cmd \times 'state and s :: 'state and S*

defines *S cf bT $\equiv \exists n. \text{qsend (cf ## bT)} = n \wedge \text{eff-at cf bT n} \approx s$*

defines *Ps cf $\equiv \mathcal{P}(bT \text{ in } T.T cf. S cf bT)$*

defines *N cf n bT $\equiv \neg \text{discrCf ((cf ## bT) !! n)$*

defines *Pn cf n $\equiv \mathcal{P}(bT \text{ in } T.T cf. N cf n bT)$*

assumes *bisim: proper (fst cf1) proper (fst cf2) fst cf1 ≈ 01 fst cf2 snd cf1 \approx snd cf2*

shows *dist (Ps cf1) (Ps cf2) $\leq Pn cf1 n + Pn cf2 m$*

$\langle \text{proof} \rangle$

lemma *AE-T-max-qsend-time*:

fixes *cf* **and** *e* :: *real* **assumes** *AE*: *AE bT in T.T cf. qsend (cf ## bT) < ∞*
0 < e
shows $\exists N. \mathcal{P}(bT \text{ in } T.T \text{ cf. } \neg \text{discrCf } ((cf \ ## \ bT) \ !! \ N)) < e$
<proof>

lemma *Ps-eq*:

fixes *cf1 cf2 s* **and** *S*
defines *S cf bT* $\equiv \exists n. \text{qsend } (cf \ ## \ bT) = n \wedge \text{eff-at } cf \ bT \ n \approx s$
defines *Ps cf* $\equiv \mathcal{P}(bT \text{ in } T.T \text{ cf. } S \text{ cf } bT)$
assumes *qsterm1*: *AE bT in T.T cf1. qsend (cf1 ## bT) < ∞*
assumes *qsterm2*: *AE bT in T.T cf2. qsend (cf2 ## bT) < ∞*
and *bisim*: *proper (fst cf1) proper (fst cf2) fst cf1 ≈01 fst cf2 snd cf1 ≈ snd cf2*
shows *Ps cf1 = Ps cf2*
<proof>

lemma *siso-trace*:

assumes *siso c s ≈ t enabled (c, t) cfT*
shows *siso (cont-at (c, s) cfT n)*
and *cont-at (c, s) cfT n = cont-at (c, t) cfT n*
and *eff-at (c, s) cfT n ≈ eff-at (c, t) cfT n*
<proof>

lemma *Sbis-trace*:

assumes *proper (fst cf1) proper (fst cf2) fst cf1 ≈s fst cf2 snd cf1 ≈ snd cf2*
shows $\mathcal{P}(cfT \text{ in } T.T \text{ cf1. eff-at } cf1 \ cfT \ n \approx s') = \mathcal{P}(cfT \text{ in } T.T \text{ cf2. eff-at } cf2 \ cfT \ n \approx s')$
(is ?P cf1 n = ?P cf2 n)
<proof>

3.4 Final Theorems

theorem *ZObis-eSec*: $\llbracket \text{proper } c; c \approx 01 \ c; \text{aeT } c \rrbracket \implies \text{eSec } c$
<proof>

theorem *Sbis-amSec*: $\llbracket \text{proper } c; c \approx s \ c \rrbracket \implies \text{amSec } c$
<proof>

theorem *amSec-eSec*:

assumes [*simp*]: *proper c* **and** *aeT c* **amSec c** **shows** *eSec c*
<proof>

end

end

4 Compositionality of Resumption-Based Noninterference

theory *Compositionality*
imports *Resumption-Based*
begin

context *PL-Indis*
begin

4.1 Compatibility and discreteness of atoms, tests and choices

definition *compatAtm* **where**
compatAtm atm \equiv
ALL s t. s \approx t \longrightarrow aval atm s \approx aval atm t

definition *presAtm* **where**
presAtm atm \equiv
ALL s. s \approx aval atm s

definition *compatTst* **where**
compatTst tst \equiv
ALL s t. s \approx t \longrightarrow tval tst s = tval tst t

lemma *discrAt-compatAt[simp]*:
assumes *presAtm atm*
shows *compatAtm atm*
<proof>

definition *compatCh* **where**
compatCh ch $\equiv \forall s t. s \approx t \longrightarrow cval ch s = cval ch t$

lemma *compatCh-cval[simp]*:
assumes *compatCh ch* **and** *s \approx t*
shows *cval ch s = cval ch t*
<proof>

4.2 Compositionality of self-isomorphism

Self-Isomorphism versus language constructs:

lemma *siso-Done[simp]*:
siso Done
<proof>

lemma *siso-Atm[simp]*:
siso (Atm atm) = compatAtm atm
<proof>

lemma *siso-Seq*[simp]:
assumes *: *siso c1* **and** **: *siso c2*
shows *siso (c1 ;; c2)*
⟨*proof*⟩

lemma *siso-While*[simp]:
assumes *compatTst tst* **and** *siso c*
shows *siso (While tst c)*
⟨*proof*⟩

lemma *siso-Ch*[simp]:
assumes *compatCh ch*
and *: *siso c1* **and** **: *siso c2*
shows *siso (Ch ch c1 c2)*
⟨*proof*⟩

lemma *siso-Par*[simp]:
assumes *properL cl* **and** *sisoL cl*
shows *siso (Par cl)*
⟨*proof*⟩

lemma *siso-ParT*[simp]:
assumes *properL cl* **and** *sisoL cl*
shows *siso (ParT cl)*
⟨*proof*⟩

Self-isomorphism implies strong bisimilarity:

lemma *bij-betw-emp*[simp]:
bij-betw f {} {}
⟨*proof*⟩

lemma *part-full*[simp]:
part I {I}
⟨*proof*⟩

definition *singlPart* **where**
singlPart I $\equiv \{\{i\} \mid i . i \in I\}$

lemma *part-singlPart*[simp]:
part I (singlPart I)
⟨*proof*⟩

lemma *singlPart-inj-on*[simp]:
inj-on (image f) (singlPart I) = inj-on f I
⟨*proof*⟩

lemma *singlPart-surj*[simp]:
(image f) ` (singlPart I) = (singlPart J) \longleftrightarrow f ` I = J

$\langle proof \rangle$

lemma *singlPart-bij-betw*[simp]:

$bij\text{-betw } (image\ f) (singlPart\ I) (singlPart\ J) = bij\text{-betw } f\ I\ J$

$\langle proof \rangle$

lemma *singlPart-finite1*:

assumes *finite* (*singlPart I*)

shows *finite* (*I::'a set*)

$\langle proof \rangle$

lemma *singlPart-finite*[simp]:

$finite\ (singlPart\ I) = finite\ I$

$\langle proof \rangle$

lemma *emp-notIn-singlPart*[simp]:

$\{\} \notin singlPart\ I$

$\langle proof \rangle$

lemma *Sbis-coinduct*[*consumes 1, case-names step, coinduct set*]:

$R\ c\ d \implies$

$(\bigwedge c\ d\ s\ t. R\ c\ d \implies s \approx t \implies$

$\exists P\ F. mC\text{-}C\text{-part}\ c\ d\ P\ F \wedge inj\text{-on}\ F\ P \wedge mC\text{-}C\text{-wt}\ c\ d\ s\ t\ P\ F \wedge$

$(\forall I \in P. \forall i \in I. \forall j \in F\ I.$

$eff\ c\ s\ i \approx eff\ d\ t\ j \wedge (R\ (cont\ c\ s\ i)\ (cont\ d\ t\ j) \vee (cont\ c\ s\ i, cont\ d$

$t\ j) \in Sbis)))$

$\implies (c, d) \in Sbis$

$\langle proof \rangle$

lemma *siso-Sbis*[simp]: $siso\ c \implies c \approx s\ c$

$\langle proof \rangle$

4.3 Discreetness versus language constructs:

lemma *discr-Done*[simp]: $discr\ Done$

$\langle proof \rangle$

lemma *discr-Atm-presAtm*[simp]: $discr\ (Atm\ atm) = presAtm\ atm$

$\langle proof \rangle$

lemma *discr-Seq*[simp]:

$discr\ c1 \implies discr\ c2 \implies discr\ (c1\ ;;\ c2)$

$\langle proof \rangle$

lemma *discr-While*[simp]: **assumes** $discr\ c$ **shows** $discr\ (While\ tst\ c)$

$\langle proof \rangle$

lemma *discr-Ch*[simp]: $discr\ c1 \implies discr\ c2 \implies discr\ (Ch\ ch\ c1\ c2)$

$\langle proof \rangle$

lemma *discr-Par[simp]*: $\text{properL } cl \implies \text{discrL } cl \implies \text{discr } (\text{Par } cl)$
(proof)

lemma *discr-ParT[simp]*: $\text{properL } cl \implies \text{discrL } cl \implies \text{discr } (\text{ParT } cl)$
(proof)

lemma *discr-finished[simp]*: $\text{proper } c \implies \text{finished } c \implies \text{discr } c$
(proof)

4.4 Strong bisimilarity versus language constructs

lemma *Sbis-pres-discr-L*:
 $c \approx_s d \implies \text{discr } d \implies \text{discr } c$
(proof)

lemma *Sbis-pres-discr-R*:
assumes *discr c and c ≈_s d*
shows *discr d*
(proof)

lemma *Sbis-finished-discr-L*:
assumes *c ≈_s d and proper d and finished d*
shows *discr c*
(proof)

lemma *Sbis-finished-discr-R*:
assumes *proper c and finished c and c ≈_s d*
shows *discr d*
(proof)

definition *thetaSD* **where**
 $\text{thetaSD} \equiv \{(c, d) \mid c \ d . \text{proper } c \wedge \text{proper } d \wedge \text{discr } c \wedge \text{discr } d\}$

lemma *thetaSD-Sretr*:
 $\text{thetaSD} \subseteq \text{Sretr } \text{thetaSD}$
(proof)

lemma *thetaSD-Sbis*:
 $\text{thetaSD} \subseteq \text{Sbis}$
(proof)

theorem *discr-Sbis[simp]*:
assumes *proper c and proper d and discr c and discr d*
shows *c ≈_s d*
(proof)

definition *thetaSDone* **where**
thetaSDone $\equiv \{(Done, Done)\}$

lemma *thetaSDone-Sretr*:
thetaSDone \subseteq *Sretr thetaSDone*
(proof)

lemma *thetaSDone-Sbis*:
thetaSDone \subseteq *Sbis*
(proof)

theorem *Done-Sbis[simp]*:
Done \approx_s *Done*
(proof)

definition *thetaSAtm* **where**
thetaSAtm atm \equiv
 $\{(Atm\ atm, Atm\ atm), (Done, Done)\}$

lemma *thetaSAtm-Sretr*:
assumes *compatAtm atm*
shows *thetaSAtm atm* \subseteq *Sretr (thetaSAtm atm)*
(proof)

lemma *thetaSAtm-Sbis*:
assumes *compatAtm atm*
shows *thetaSAtm atm* \subseteq *Sbis*
(proof)

theorem *Atm-Sbis[simp]*:
assumes *compatAtm atm*
shows *Atm atm* \approx_s *Atm atm*
(proof)

definition *thetaSSeqI* **where**
thetaSSeqI \equiv
 $\{(e ;; c, e ;; d) \mid e\ c\ d .\ siso\ e \wedge c \approx_s d\}$

lemma *thetaSSeqI-Sretr*:
thetaSSeqI \subseteq *Sretr (thetaSSeqI Un Sbis)*
(proof)

lemma *thetaSSeqI-Sbis*:
thetaSSeqI \subseteq *Sbis*
(proof)

theorem *Seq-iso-Sbis*[simp]:
assumes *siso e and c2 ≈s d2*
shows *e ;; c2 ≈s e ;; d2*
⟨proof⟩

definition *thetaSSeqD* where
thetaSSeqD ≡
{(*c1* ;; *c2*, *d1* ;; *d2*) |
c1 c2 d1 d2.
proper c1 ∧ proper d1 ∧ proper c2 ∧ proper d2 ∧
discr c2 ∧ discr d2 ∧
c1 ≈s d1}

lemma *thetaSSeqD-Sretr*:
thetaSSeqD ⊆ *Sretr (thetaSSeqD Un Sbis)*
⟨proof⟩

lemma *thetaSSeqD-Sbis*:
thetaSSeqD ⊆ *Sbis*
⟨proof⟩

theorem *Seq-Sbis*[simp]:
assumes *proper c1 and proper d1 and proper c2 and proper d2*
and *c1 ≈s d1 and discr c2 and discr d2*
shows *c1 ;; c2 ≈s d1 ;; d2*
⟨proof⟩

definition *thetaSCh* where
thetaSCh ch c1 c2 d1 d2 ≡ {(*Ch ch c1 c2*, *Ch ch d1 d2*)}

lemma *thetaSCh-Sretr*:
assumes *compatCh ch and c1 ≈s d1 and c2 ≈s d2*
shows *thetaSCh ch c1 c2 d1 d2* ⊆
Sretr (thetaSCh ch c1 c2 d1 d2 ∪ Sbis)
(is ?th ⊆ *Sretr (?th ∪ Sbis)*)
⟨proof⟩

lemma *thetaSCh-Sbis*:
assumes *compatCh ch and c1 ≈s d1 and c2 ≈s d2*
shows *thetaSCh ch c1 c2 d1 d2* ⊆ *Sbis*
⟨proof⟩

theorem *Ch-iso-Sbis*[simp]:
assumes *compatCh ch and c1 ≈s d1 and c2 ≈s d2*

shows $Ch\ ch\ c1\ c2 \approx_s Ch\ ch\ d1\ d2$
(proof)

definition *shift* **where**
 $shift\ cl\ n \equiv image\ (\%i.\ brnL\ cl\ n + i)$

definition *back* **where**
 $back\ cl\ n \equiv image\ (\%ii.\ ii - brnL\ cl\ n)$

lemma *emp-shift*[simp]:
 $shift\ cl\ n\ I = \{\} \longleftrightarrow I = \{\}$
(proof)

lemma *emp-shift-rev*[simp]:
 $\{\} = shift\ cl\ n\ I \longleftrightarrow I = \{\}$
(proof)

lemma *emp-back*[simp]:
 $back\ cl\ n\ II = \{\} \longleftrightarrow II = \{\}$
(proof)

lemma *emp-back-rev*[simp]:
 $\{\} = back\ cl\ n\ II \longleftrightarrow II = \{\}$
(proof)

lemma *in-shift*[simp]:
 $brnL\ cl\ n + i \in shift\ cl\ n\ I \longleftrightarrow i \in I$
(proof)

lemma *in-back*[simp]:
 $ii \in II \implies ii - brnL\ cl\ n \in back\ cl\ n\ II$
(proof)

lemma *in-back2*[simp]:
assumes $ii > brnL\ cl\ n$ **and** $II \subseteq \{brnL\ cl\ n ..<+ brn\ (cl!n)\}$
shows $ii - brnL\ cl\ n \in back\ cl\ n\ II \longleftrightarrow ii \in II$ (**is** ?L \longleftrightarrow ?R)
(proof)

lemma *shift*[simp]:
assumes $I \subseteq \{..< brn\ (cl!n)\}$
shows $shift\ cl\ n\ I \subseteq \{brnL\ cl\ n ..<+ brn\ (cl!n)\}$
(proof)

lemma *shift2*[simp]:
assumes $I \subseteq \{..< brn\ (cl!n)\}$
and $ii \in shift\ cl\ n\ I$
shows $brnL\ cl\ n \leq ii \wedge ii < brnL\ cl\ n + brn\ (cl!n)$

<proof>

lemma *shift3[simp]*:

assumes $n: n < \text{length } cl$ **and** $I: I \subseteq \{..
n\}$

and $ii: ii \in \text{shift } cl \ n \ I$

shows $ii < \text{brnL } cl \ (\text{length } cl)$

<proof>

lemma *back[simp]*:

assumes $II \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$

shows $\text{back } cl \ n \ II \subseteq \{..
n\}$

<proof>

lemma *back2[simp]*:

assumes $II \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$

and $i \in \text{back } cl \ n \ II$

shows $i < \text{brn } (cl!n)$

<proof>

lemma *shift-inj[simp]*:

$\text{shift } cl \ n \ I1 = \text{shift } cl \ n \ I2 \longleftrightarrow I1 = I2$

<proof>

lemma *shift-mono[simp]*:

$\text{shift } cl \ n \ I1 \subseteq \text{shift } cl \ n \ I2 \longleftrightarrow I1 \subseteq I2$

<proof>

lemma *shift-Int[simp]*:

$\text{shift } cl \ n \ I1 \cap \text{shift } cl \ n \ I2 = \{\} \longleftrightarrow I1 \cap I2 = \{\}$

<proof>

lemma *inj-shift*: $\text{inj } (\text{shift } cl \ n)$

<proof>

lemma *inj-on-shift*: $\text{inj-on } (\text{shift } cl \ n) \ K$

<proof>

lemma *back-shift[simp]*:

$\text{back } cl \ n \ (\text{shift } cl \ n \ I) = I$

<proof>

lemma *shift-back[simp]*:

assumes $II \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$

shows $\text{shift } cl \ n \ (\text{back } cl \ n \ II) = II$

<proof>

lemma *back-inj[simp]*:

assumes $II1: II1 \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$

and $II2: II2 \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$

shows $\text{back } cl \ n \ II1 = \text{back } cl \ n \ II2 \longleftrightarrow II1 = II2$ (is $?L = ?R \longleftrightarrow II1 = II2$)
 ⟨proof⟩

lemma *back-mono*[simp]:
assumes $II1 \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$
and $II2 \subseteq \{\text{brnL } cl \ n \ ..< \text{brnL } cl \ n + \text{brn } (cl!n)\}$
shows $\text{back } cl \ n \ II1 \subseteq \text{back } cl \ n \ II2 \longleftrightarrow II1 \subseteq II2$
 (is $?L \subseteq ?R \longleftrightarrow II1 \subseteq II2$)
 ⟨proof⟩

lemma *back-Int*[simp]:
assumes $II1 \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$
and $II2 \subseteq \{\text{brnL } cl \ n \ ..< \text{brnL } cl \ n + \text{brn } (cl!n)\}$
shows $\text{back } cl \ n \ II1 \cap \text{back } cl \ n \ II2 = \{\} \longleftrightarrow II1 \cap II2 = \{\}$
 (is $?L \cap ?R = \{\} \longleftrightarrow II1 \cap II2 = \{\}$)
 ⟨proof⟩

lemma *inj-on-back*:
inj-on ($\text{back } cl \ n$) ($\text{Pow } \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$)
 ⟨proof⟩

lemma *shift-surj*:
assumes $II \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\}$
shows $\exists I. I \subseteq \{\text{brnL } cl \ n \ ..< \text{brn } (cl!n)\} \wedge \text{shift } cl \ n \ I = II$
 ⟨proof⟩

lemma *back-surj*:
assumes $I \subseteq \{\text{brnL } cl \ n \ ..< \text{brn } (cl!n)\}$
shows $\exists II. II \subseteq \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\} \wedge \text{back } cl \ n \ II = I$
 ⟨proof⟩

lemma *shift-part*[simp]:
assumes $\text{part } \{\text{brnL } cl \ n \ ..< \text{brn } (cl!n)\} \ P$
shows $\text{part } \{\text{brnL } cl \ n \ ..<+ \text{brn } (cl!n)\} \ (\text{shift } cl \ n \ ' P)$
 ⟨proof⟩

lemma *part-brn-disj1*:
assumes $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{\text{brnL } cl \ n \ ..< \text{brn } (cl!n)\} \ (P \ n)$
and $n1: n1 < \text{length } cl$ **and** $n2: n2 < \text{length } cl$
and $II1: II1 \in \text{shift } cl \ n1 \ ' (P \ n1)$ **and** $II2: II2 \in \text{shift } cl \ n2 \ ' (P \ n2)$ **and** $d: n1 \neq n2$
shows $II1 \cap II2 = \{\}$
 ⟨proof⟩

lemma *part-brn-disj2*:
assumes $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{\text{brnL } cl \ n \ ..< \text{brn } (cl!n)\} \ (P \ n) \wedge \{\} \notin P \ n$
and $n1: n1 < \text{length } cl$ **and** $n2: n2 < \text{length } cl$ **and** $d: n1 \neq n2$
shows $\text{shift } cl \ n1 \ ' (P \ n1) \cap \text{shift } cl \ n2 \ ' (P \ n2) = \{\}$ (is $?L \cap ?R = \{\}$)
 ⟨proof⟩

lemma *part-brn-disj3*:
assumes $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n)$
and $n1: n1 < \text{length } cl$ **and** $n2: n2 < \text{length } cl$
and $I1: I1 \in P \ n1$ **and** $I2: I2 \in P \ n2$ **and** $d: n1 \neq n2$
shows $\text{shift } cl \ n1 \ I1 \cap \text{shift } cl \ n2 \ I2 = \{\}$
 $\langle \text{proof} \rangle$

lemma *sum-wt-Par-sub-shift[simp]*:
assumes $cl: \text{properL } cl$ **and** $n: n < \text{length } cl$ **and**
 $I: I \subseteq \{.. < \text{brn } (cl!n)\}$
shows
 $\text{sum } (wt \ (Par \ cl) \ s) \ (\text{shift } cl \ n \ I) =$
 $1 / (\text{length } cl) * \text{sum } (wt \ (cl!n) \ s) \ I$
 $\langle \text{proof} \rangle$

lemma *sum-wt-ParT-sub-WtFT-pickFT-0-shift[simp]*:
assumes $cl: \text{properL } cl$ **and** $nf: WtFT \ cl = 1$
and $I: I \subseteq \{.. < \text{brn } (cl! (\text{pickFT } cl))\}$ $0 \in I$
shows
 $\text{sum } (wt \ (ParT \ cl) \ s) \ (\text{shift } cl \ (\text{pickFT } cl) \ I) = 1$
 $\langle \text{proof} \rangle$

lemma *sum-wt-ParT-sub-WtFT-notPickFT-0-shift[simp]*:
assumes $cl: \text{properL } cl$ **and** $nf: WtFT \ cl = 1$ **and** $n: n < \text{length } cl$
and $I: I \subseteq \{.. < \text{brn } (cl!n)\}$ **and** $nI: n = \text{pickFT } cl \longrightarrow 0 \notin I$
shows $\text{sum } (wt \ (ParT \ cl) \ s) \ (\text{shift } cl \ n \ I) = 0$
 $\langle \text{proof} \rangle$

lemma *sum-wt-ParT-sub-notWtFT-finished-shift[simp]*:
assumes $cl: \text{properL } cl$ **and** $nf: WtFT \ cl \neq 1$ **and** $n: n < \text{length } cl$ **and** $cln:$
 $\text{finished } (cl!n)$
and $I: I \subseteq \{.. < \text{brn } (cl!n)\}$
shows $\text{sum } (wt \ (ParT \ cl) \ s) \ (\text{shift } cl \ n \ I) = 0$
 $\langle \text{proof} \rangle$

lemma *sum-wt-ParT-sub-notWtFT-notFinished-shift[simp]*:
assumes $cl: \text{properL } cl$ **and** $nf: WtFT \ cl \neq 1$
and $n: n < \text{length } cl$ **and** $cln: \neg \text{finished } (cl!n)$
and $I: I \subseteq \{.. < \text{brn } (cl!n)\}$
shows
 $\text{sum } (wt \ (ParT \ cl) \ s) \ (\text{shift } cl \ n \ I) =$
 $(1 / (\text{length } cl)) / (1 - WtFT \ cl) * \text{sum } (wt \ (cl!n) \ s) \ I$
 $\langle \text{proof} \rangle$

definition *UNpart where*
 $UNpart \ cl \ P \equiv \bigcup_{n < \text{length } cl.} \text{shift } cl \ n \ ' (P \ n)$

lemma *UNpart-cases*[*elim, consumes 1, case-names Local*]:
assumes $II \in \text{UNpart } cl \ P$ **and**
 $\bigwedge n \ I. \llbracket n < \text{length } cl; I \in P \ n; II = \text{shift } cl \ n \ I \rrbracket \implies phi$
shows *phi*
<proof>

lemma *emp-UNpart*:
assumes $\bigwedge n. n < \text{length } cl \implies \{\} \notin P \ n$
shows $\{\} \notin \text{UNpart } cl \ P$
<proof>

lemma *part-UNpart*:
assumes *cl*: *properL cl* **and**
 $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n)$
shows $\text{part } \{.. < \text{brnL } cl \ (\text{length } cl)\} (\text{UNpart } cl \ P)$
(is part ?J ?Q)
<proof>

definition *pickT-pred* **where**
 $\text{pickT-pred } cl \ P \ II \ n \equiv n < \text{length } cl \wedge II \in \text{shift } cl \ n \ ' (P \ n)$

definition *pickT* **where**
 $\text{pickT } cl \ P \ II \equiv \text{SOME } n. \text{pickT-pred } cl \ P \ II \ n$

lemma *pickT-pred*:
assumes $II \in \text{UNpart } cl \ P$
shows $\exists n. \text{pickT-pred } cl \ P \ II \ n$
<proof>

lemma *pickT-pred-unique*:
assumes $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$
and $1: \text{pickT-pred } cl \ P \ II \ n1$ **and** $2: \text{pickT-pred } cl \ P \ II \ n2$
shows $n1 = n2$
<proof>

lemma *pickT-pred-pickT*:
assumes $II \in \text{UNpart } cl \ P$
shows $\text{pickT-pred } cl \ P \ II \ (\text{pickT } cl \ P \ II)$
<proof>

lemma *pickT-pred-pickT-unique*:
assumes $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$
and $\text{pickT-pred } cl \ P \ II \ n$
shows $n = \text{pickT } cl \ P \ II$
<proof>

lemma *pickT-length[simp]*:
assumes $II \in UNpart\ cl\ P$
shows $pickT\ cl\ P\ II < length\ cl$
 $\langle proof \rangle$

lemma *pickT-shift[simp]*:
assumes $II \in UNpart\ cl\ P$
shows $II \in shift\ cl\ (pickT\ cl\ P\ II) \text{ ' } (P\ (pickT\ cl\ P\ II))$
 $\langle proof \rangle$

lemma *pickT-unique*:
assumes $P: \bigwedge n. n < length\ cl \implies part\ \{..< brn\ (cl!n)\}\ (P\ n) \wedge \{\} \notin P\ n$
and $n < length\ cl$ **and** $II \in shift\ cl\ n \text{ ' } (P\ n)$
shows $n = pickT\ cl\ P\ II$
 $\langle proof \rangle$

definition *UNlift where*
 $UNlift\ cl\ dl\ P\ F\ II \equiv$
 $shift\ dl\ (pickT\ cl\ P\ II)\ (F\ (pickT\ cl\ P\ II)\ (back\ cl\ (pickT\ cl\ P\ II)\ II))$

lemma *UNlift-shift[simp]*:
assumes $P: \bigwedge n. n < length\ cl \implies part\ \{..< brn\ (cl!n)\}\ (P\ n) \wedge \{\} \notin P\ n$
and $n: n < length\ cl$ **and** $I: I \in P\ n$
shows $UNlift\ cl\ dl\ P\ F\ (shift\ cl\ n\ I) = shift\ dl\ n\ (F\ n\ I)$
 $\langle proof \rangle$

lemma *UNlift-inj-on*:
assumes $l: length\ cl = length\ dl$
and $P: \bigwedge n. n < length\ cl \implies part\ \{..< brn\ (cl!n)\}\ (P\ n) \wedge \{\} \notin P\ n$
and $FP: \bigwedge n. n < length\ dl \implies part\ \{..< brn\ (dl!n)\}\ (F\ n \text{ ' } (P\ n)) \wedge \{\} \notin F\ n \text{ ' } (P\ n)$
and $F: \bigwedge n. n < length\ cl \implies inj\text{-on}\ (F\ n)\ (P\ n)$
shows $inj\text{-on}\ (UNlift\ cl\ dl\ P\ F)\ (UNpart\ cl\ P)\ (\mathbf{is}\ inj\text{-on}\ ?G\ ?Q)$
 $\langle proof \rangle$

lemma *UNlift-UNpart*:
assumes $l: length\ cl = length\ dl$
and $P: \bigwedge n. n < length\ cl \implies part\ \{..< brn\ (cl!n)\}\ (P\ n) \wedge \{\} \notin P\ n$
shows $(UNlift\ cl\ dl\ P\ F) \text{ ' } (UNpart\ cl\ P) = UNpart\ dl\ (\%n. F\ n \text{ ' } (P\ n))\ (\mathbf{is}\ ?G \text{ ' } ?Q = ?R)$
 $\langle proof \rangle$

lemma *emp-UNlift-UNpart*:
assumes $l: length\ cl = length\ dl$
and $P: \bigwedge n. n < length\ cl \implies part\ \{..< brn\ (cl!n)\}\ (P\ n) \wedge \{\} \notin P\ n$
and $FP: \bigwedge n. n < length\ dl \implies \{\} \notin F\ n \text{ ' } (P\ n)$
shows $\{\} \notin (UNlift\ cl\ dl\ P\ F) \text{ ' } (UNpart\ cl\ P)\ (\mathbf{is}\ \{\} \notin ?R)$
 $\langle proof \rangle$

lemma *part-UNlift-UNpart*:

assumes l : $\text{length } cl = \text{length } dl$ **and** dl : $\text{properL } dl$

and P : $\bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

and FP : $\bigwedge n. n < \text{length } dl \implies \text{part } \{.. < \text{brn } (dl!n)\} (F \ n \ ' (P \ n))$

shows $\text{part } \{.. < \text{brnL } dl \ (\text{length } dl)\} ((\text{UNlift } cl \ dl \ P \ F) \ ' (\text{UNpart } cl \ P))$ **(is part**

$?C \ ?R$)

$\langle \text{proof} \rangle$

lemma *ss-wt-Par-UNlift*:

assumes l : $\text{length } cl = \text{length } dl$

and $cldl$: $\text{properL } cl \ \text{properL } dl$ **and** II : $II \in \text{UNpart } cl \ P$

and P : $\bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

and FP : $\bigwedge n. n < \text{length } dl \implies \text{part } \{.. < \text{brn } (dl!n)\} (F \ n \ ' (P \ n))$

and sw :

$\bigwedge n \ I. \llbracket n < \text{length } cl; I \in P \ n \rrbracket \implies$

$\text{sum } (wt \ (cl \ ! \ n) \ s) \ I =$

$\text{sum } (wt \ (dl \ ! \ n) \ t) \ (F \ n \ I)$

and st : $s \approx t$

shows

$\text{sum } (wt \ (Par \ cl) \ s) \ II =$

$\text{sum } (wt \ (Par \ dl) \ t) \ (\text{UNlift } cl \ dl \ P \ F \ II)$ **(is ?L = ?R)**

$\langle \text{proof} \rangle$

definition *thetaSPar* **where**

$\text{thetaSPar} \equiv$

$\{(Par \ cl, \ Par \ dl) \mid$
 $\quad cl \ dl. \ \text{properL } cl \wedge \text{properL } dl \wedge \text{SbisL } cl \ dl\}$

lemma *cont-eff-Par-UNlift*:

assumes l : $\text{length } cl = \text{length } dl$

and $cldl$: $\text{properL } cl \ \text{properL } dl \ \text{SbisL } cl \ dl$

and II : $II \in \text{UNpart } cl \ P$ **and** ii : $ii \in II$ **and** jj : $jj \in \text{UNlift } cl \ dl \ P \ F \ II$

and P : $\bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

and FP : $\bigwedge n. n < \text{length } dl \implies \text{part } \{.. < \text{brn } (dl!n)\} (F \ n \ ' (P \ n))$

and eff-cont :

$\bigwedge n \ I \ i \ j. \llbracket n < \text{length } cl; I \in P \ n; i \in I; j \in F \ n \ I \rrbracket \implies$

$\text{eff } (cl!n) \ s \ i \approx \text{eff } (dl!n) \ t \ j \wedge$

$\text{cont } (cl!n) \ s \ i \approx_s \text{cont } (dl!n) \ t \ j$

and st : $s \approx t$

shows

$\text{eff } (Par \ cl) \ s \ ii \approx \text{eff } (Par \ dl) \ t \ jj \wedge$

$(\text{cont } (Par \ cl) \ s \ ii, \ \text{cont } (Par \ dl) \ t \ jj) \in \text{thetaSPar}$

(is ?eff \wedge ?cont)

$\langle \text{proof} \rangle$

lemma *thetaSPar-Sretr*: $\text{thetaSPar} \subseteq \text{Sretr } (\text{thetaSPar})$

$\langle \text{proof} \rangle$

lemma *thetaSPar-Sbis*: $\text{thetaSPar} \subseteq \text{Sbis}$
(proof)

theorem *Par-Sbis[simp]*:
assumes *properL cl* **and** *properL dl SbisL cl dl*
shows $\text{Par } cl \approx_s \text{Par } dl$
(proof)

4.5 01-bisimilarity versus language constructs

lemma *ZObis-pres-discr-L*: $c \approx_{01} d \implies \text{discr } d \implies \text{discr } c$
(proof)

theorem *ZObis-pres-discr-R*:
assumes *discr c* **and** $c \approx_{01} d$
shows *discr d*
(proof)

theorem *ZObis-finished-discr-L*:
assumes $c \approx_{01} d$ **and** *proper d* **and** *finished d*
shows *discr c*
(proof)

theorem *ZObis-finished-discr-R*:
assumes *proper c* **and** *finished c* **and** $c \approx_{01} d$
shows *discr d*
(proof)

theorem *discr-ZObis[simp]*:
assumes *proper c* **and** *proper d* **and** *discr c* **and** *discr d*
shows $c \approx_{01} d$
(proof)

theorem *Done-ZObis[simp]*:
Done \approx_{01} *Done*
(proof)

theorem *Atm-ZObis[simp]*:
assumes *compatAtm atm*
shows $\text{Atm } atm \approx_{01} \text{Atm } atm$
(proof)

definition *thetaZOSeqI* **where**
thetaZOSeqI \equiv

$\{(e ;; c, e ;; d) \mid e c d . \text{siso } e \wedge c \approx 01 d\}$

lemma *thetaZOSeqI-ZOretr*:
thetaZOSeqI \subseteq *ZOretr* (*thetaZOSeqI* *Un* *ZObis*)
(*proof*)

lemma *thetaZOSeqI-ZObis*:
thetaZOSeqI \subseteq *ZObis*
(*proof*)

theorem *Seq-siso-ZObis[simp]*:
assumes *siso e* **and** *c2 ≈ 01 d2*
shows *e ;; c2 ≈ 01 e ;; d2*
(*proof*)

definition *thetaZOSeqD* **where**
thetaZOSeqD \equiv
 $\{(c1 ;; c2, d1 ;; d2) \mid$
 c1 c2 d1 d2.
 proper c1 \wedge proper d1 \wedge proper c2 \wedge proper d2 \wedge
 discr c2 \wedge discr d2 \wedge
 c1 ≈ 01 d1\}

lemma *thetaZOSeqD-ZOretr*:
thetaZOSeqD \subseteq *ZOretr* (*thetaZOSeqD* *Un* *ZObis*)
(*proof*)

lemma *thetaZOSeqD-ZObis*:
thetaZOSeqD \subseteq *ZObis*
(*proof*)

theorem *Seq-ZObis[simp]*:
assumes *proper c1* **and** *proper d1* **and** *proper c2* **and** *proper d2*
and *c1 ≈ 01 d1* **and** *discr c2* **and** *discr d2*
shows *c1 ;; c2 ≈ 01 d1 ;; d2*
(*proof*)

definition *thetaZOCh* **where**
thetaZOCh *ch c1 c2 d1 d2* $\equiv \{(Ch \text{ ch } c1 c2, Ch \text{ ch } d1 d2)\}$

lemma *thetaZOCh-Sretr*:
assumes *compatCh ch* **and** *c1 ≈ 01 d1* **and** *c2 ≈ 01 d2*
shows *thetaZOCh* *ch c1 c2 d1 d2* \subseteq
 Sretr (*thetaZOCh* *ch c1 c2 d1 d2* \cup *ZObis*)
(**is** *?th* \subseteq *Sretr* (*?th* \cup *ZObis*))
(*proof*)

lemma *thetaZOCh-ZOretr*:
assumes *compatCh ch* **and** $c1 \approx 01 d1$ **and** $c2 \approx 01 d2$
shows $\text{thetaZOCh } ch \ c1 \ c2 \ d1 \ d2 \subseteq$
 $\text{ZOretr } (\text{thetaZOCh } ch \ c1 \ c2 \ d1 \ d2 \cup \text{ZObis})$
 $\langle \text{proof} \rangle$

lemma *thetaZOCh-ZObis*:
assumes *compatCh ch* **and** $c1 \approx 01 d1$ **and** $c2 \approx 01 d2$
shows $\text{thetaZOCh } ch \ c1 \ c2 \ d1 \ d2 \subseteq \text{ZObis}$
 $\langle \text{proof} \rangle$

theorem *Ch-iso-ZObis[simp]*:
assumes *compatCh ch* **and** $c1 \approx 01 d1$ **and** $c2 \approx 01 d2$
shows $\text{Ch } ch \ c1 \ c2 \approx 01 \ \text{Ch } ch \ d1 \ d2$
 $\langle \text{proof} \rangle$

definition *theFTOne* **where**
 $\text{theFTOne } cl \ dl \equiv \text{theFT } cl \cup \text{theFT } dl$

definition *theNFTBoth* **where**
 $\text{theNFTBoth } cl \ dl \equiv \text{theNFT } cl \cap \text{theNFT } dl$

lemma *theFTOne-sym*: $\text{theFTOne } cl \ dl = \text{theFTOne } dl \ cl$
 $\langle \text{proof} \rangle$

lemma *finite-theFTOne[simp]*:
 $\text{finite } (\text{theFTOne } cl \ dl)$
 $\langle \text{proof} \rangle$

lemma *theFTOne-length-finished[simp]*:
assumes $n \in \text{theFTOne } cl \ dl$
shows $(n < \text{length } cl \wedge \text{finished } (cl!n)) \vee (n < \text{length } dl \wedge \text{finished } (dl!n))$
 $\langle \text{proof} \rangle$

lemma *theFTOne-length[simp]*:
assumes $\text{length } cl = \text{length } dl$ **and** $n \in \text{theFTOne } cl \ dl$
shows $n < \text{length } cl$ **and** $n < \text{length } dl$
 $\langle \text{proof} \rangle$

lemma *theFTOne-intro[intro]*:
assumes $\bigwedge n. (n < \text{length } cl \wedge \text{finished } (cl!n)) \vee (n < \text{length } dl \wedge \text{finished } (dl!n))$
shows $n \in \text{theFTOne } cl \ dl$
 $\langle \text{proof} \rangle$

lemma *pickFT-theFTOne[simp]*:
assumes $\text{WtFT } cl = 1$
shows $\text{pickFT } cl \in \text{theFTOne } cl \ dl$

<proof>

lemma *finite-theNFTBoth[simp]*:
finite (theNFTBoth cl dl)
<proof>

lemma *theNFTBoth-sym*: *theNFTBoth cl dl = theNFTBoth dl cl*
<proof>

lemma *theNFTBoth-length-finished[simp]*:
assumes $n \in \text{theNFTBoth } cl \ dl$
shows $n < \text{length } cl$ **and** $\neg \text{finished } (cl!n)$
and $n < \text{length } dl$ **and** $\neg \text{finished } (dl!n)$
<proof>

lemma *theNFTBoth-intro[intro]*:
assumes $\bigwedge n. n < \text{length } cl \wedge \neg \text{finished } (cl!n) \wedge n < \text{length } dl \wedge \neg \text{finished } (dl!n)$
shows $n \in \text{theNFTBoth } cl \ dl$
<proof>

lemma *theFTOne-Int-theNFTBoth[simp]*:
theFTOne cl dl \cap theNFTBoth cl dl = {}
and *theNFTBoth cl dl \cap theFTOne cl dl = {}*
<proof>

lemma *theFT-Un-theNFT-One-Both[simp]*:
assumes $\text{length } cl = \text{length } dl$
shows
theFTOne cl dl \cup theNFTBoth cl dl = {.. $<$ length cl} and
theNFTBoth cl dl \cup theFTOne cl dl = {.. $<$ length cl}
<proof>

lemma *in-theFTOne-theNFTBoth[simp]*:
assumes $n1 \in \text{theFTOne } cl \ dl$ **and** $n2 \in \text{theNFTBoth } cl \ dl$
shows $n1 \neq n2$ **and** $n2 \neq n1$
<proof>

definition *BrnFT where*
BrnFT cl dl $\equiv \bigcup n \in \text{theFTOne } cl \ dl. \{brnL \ cl \ n \ ..<+ \ brn \ (cl!n)\}$

definition *BrnNFT where*
BrnNFT cl dl $\equiv \bigcup n \in \text{theNFTBoth } cl \ dl. \{brnL \ cl \ n \ ..<+ \ brn \ (cl!n)\}$

lemma *BrnFT-elim[elim, consumes 1, case-names Local]*:
assumes $ii \in \text{BrnFT } cl \ dl$

and $\bigwedge n i. \llbracket n \in \text{theFTOne } cl \ dl; i < \text{brn } (cl!n); ii = \text{brnL } cl \ n + i \rrbracket \implies phi$
shows phi
 $\langle proof \rangle$

lemma $\text{finite-BrnFT}[simp]$:
 $finite (BrnFT \ cl \ dl)$
 $\langle proof \rangle$

lemma $\text{BrnFT-incl-brnL}[simp]$:
assumes $l: \text{length } cl = \text{length } dl$ **and** $cl: \text{properL } cl$
shows $\text{BrnFT } cl \ dl \subseteq \{.. < \text{brnL } cl \ (\text{length } cl)\}$ (**is** $?L \subseteq ?R$)
 $\langle proof \rangle$

lemma $\text{BrnNFT-elim}[elim, \text{consumes } 1, \text{case-names Local}]$:
assumes $ii \in \text{BrnNFT } cl \ dl$
and $\bigwedge n i. \llbracket n \in \text{theNFTBoth } cl \ dl; i < \text{brn } (cl!n); ii = \text{brnL } cl \ n + i \rrbracket \implies phi$
shows phi
 $\langle proof \rangle$

lemma $\text{finite-BrnNFT}[simp]$:
 $finite (BrnNFT \ cl \ dl)$
 $\langle proof \rangle$

lemma $\text{BrnNFT-incl-brnL}[simp]$:
assumes $cl: \text{properL } cl$
shows $\text{BrnNFT } cl \ dl \subseteq \{.. < \text{brnL } cl \ (\text{length } cl)\}$ (**is** $?L \subseteq ?R$)
 $\langle proof \rangle$

lemma $\text{BrnFT-Int-BrnNFT}[simp]$:
assumes $l: \text{length } cl = \text{length } dl$
shows
 $\text{BrnFT } cl \ dl \cap \text{BrnNFT } cl \ dl = \{\}$ (**is** $?L$)
and $\text{BrnNFT } cl \ dl \cap \text{BrnFT } cl \ dl = \{\}$ (**is** $?R$)
 $\langle proof \rangle$

lemma $\text{BrnFT-Un-BrnNFT}[simp]$:
assumes $l: \text{length } cl = \text{length } dl$ **and** $cl: \text{properL } cl$
shows $\text{BrnFT } cl \ dl \cup \text{BrnNFT } cl \ dl = \{.. < \text{brnL } cl \ (\text{length } cl)\}$ (**is** $?L1 = ?R$)
and $\text{BrnNFT } cl \ dl \cup \text{BrnFT } cl \ dl = \{.. < \text{brnL } cl \ (\text{length } cl)\}$ (**is** $?L2 = ?R$)
 $\langle proof \rangle$

lemma BrnFT-part :
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} (P \ n)$
shows $\text{BrnFT } cl \ dl = (\bigcup n \in \text{theFTOne } cl \ dl. \text{Union } (\text{shift } cl \ n \ ' (P \ n)))$ (**is** $?L = ?R$)
 $\langle proof \rangle$

lemma $\text{brnL-pickFT-BrnFT}[simp]$:

assumes $\text{properL } cl \text{ and } \text{WtFT } cl = 1$
shows $\text{brnL } cl \text{ (pickFT } cl) \in \text{BrnFT } cl \text{ dl}$
 $\langle \text{proof} \rangle$

lemma $\text{WtFT-ParT-BrnFT[simp]}$:
assumes $\text{length } cl = \text{length } dl \text{ properL } cl \text{ and } \text{WtFT } cl = 1$
shows $\text{sum } (wt \text{ (ParT } cl) s) (\text{BrnFT } cl \text{ dl}) = 1$
 $\langle \text{proof} \rangle$

definition UNpart1 **where**
 $\text{UNpart1 } cl \text{ dl } P \equiv \bigcup n \in \text{theNFTBoth } cl \text{ dl. shift } cl \text{ } n \text{ } (P \text{ } n)$

definition UNpart01 **where**
 $\text{UNpart01 } cl \text{ dl } P \equiv \{\text{BrnFT } cl \text{ dl}\} \cup \text{UNpart1 } cl \text{ dl } P$

lemma $\text{BrnFT-UNpart01[simp]}$:
 $\text{BrnFT } cl \text{ dl} \in \text{UNpart01 } cl \text{ dl } P$
 $\langle \text{proof} \rangle$

lemma $\text{UNpart1-cases[elim, consumes 1, case-names Local]}$:
assumes $II \in \text{UNpart1 } cl \text{ dl } P$
 $\bigwedge n \text{ } I. \llbracket n \in \text{theNFTBoth } cl \text{ dl}; I \in P \text{ } n; II = \text{shift } cl \text{ } n \text{ } I \rrbracket \implies \text{phi}$
shows phi
 $\langle \text{proof} \rangle$

lemma $\text{UNpart01-cases[elim, consumes 1, case-names Local0 Local]}$:
assumes $II \in \text{UNpart01 } cl \text{ dl } P \text{ and } II = \text{BrnFT } cl \text{ dl} \implies \text{phi}$
 $\bigwedge n \text{ } I. \llbracket n \in \text{theNFTBoth } cl \text{ dl}; I \in P \text{ } n; II = \text{shift } cl \text{ } n \text{ } I; II \in \text{UNpart1 } cl \text{ dl } P \rrbracket$
 $\implies \text{phi}$
shows phi
 $\langle \text{proof} \rangle$

lemma emp-UNpart1 :
assumes $\bigwedge n. n < \text{length } cl \implies \{\} \notin P \text{ } n$
shows $\{\} \notin \text{UNpart1 } cl \text{ dl } P$
 $\langle \text{proof} \rangle$

lemma emp-UNpart01 :
assumes $\bigwedge n. n < \text{length } cl \implies \{\} \notin P \text{ } n$
shows $\{\} \notin \text{UNpart01 } cl \text{ dl } P - \{\text{BrnFT } cl \text{ dl}\}$
 $\langle \text{proof} \rangle$

lemma $\text{BrnFT-Int-UNpart1[simp]}$:
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl \ll n)\} (P \text{ } n) \wedge \{\} \notin P \text{ } n$
and $II: II \in \text{UNpart1 } cl \text{ dl } P$
shows $\text{BrnFT } cl \text{ dl} \cap II = \{\}$

<proof>

lemma *BrnFT-notIn-UNpart1*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

shows $\text{BrnFT } cl \ dl \notin \text{UNpart1 } cl \ dl \ P$

<proof>

lemma *UNpart1-UNpart01*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

shows $\text{UNpart1 } cl \ dl \ P = \text{UNpart01 } cl \ dl \ P - \{\text{BrnFT } cl \ dl\}$

<proof>

lemma *part-UNpart1[simp]*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n)$

shows $\text{part } (\text{BrnNFT } cl \ dl) (\text{UNpart1 } cl \ dl \ P)$

<proof>

lemma *part-UNpart01*:

assumes *cl*: *properL cl* **and** *l*: $\text{length } cl = \text{length } dl$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

shows $\text{part } \{..< \text{brnL } cl \ (\text{length } cl)\} (\text{UNpart01 } cl \ dl \ P)$

<proof>

definition *UNlift01 where*

UNlift01 cl dl P F II \equiv

if $II = \text{BrnFT } cl \ dl$

then $\text{BrnFT } dl \ cl$

else $\text{shift } dl \ (\text{pickT } cl \ P \ II) \ (F \ (\text{pickT } cl \ P \ II) \ (\text{back } cl \ (\text{pickT } cl \ P \ II) \ II))$

lemma *UNlift01-BrnFT[simp]*:

UNlift01 cl dl P F (BrnFT cl dl) = BrnFT dl cl

<proof>

lemma *UNlift01-shift[simp]*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

and *n*: $n \in \text{theNFTBoth } cl \ dl$ **and** *I*: $I \in P \ n$

shows $\text{UNlift01 } cl \ dl \ P \ F \ (\text{shift } cl \ n \ I) = \text{shift } dl \ n \ (F \ n \ I)$

<proof>

lemma *UNlift01-inj-on-UNpart1*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$

and *FP*: $\bigwedge n. n < \text{length } dl \implies \text{part } \{..< \text{brn } (dl!n)\} (F \ n \ ' (P \ n)) \wedge \{\} \notin F \ n \ '$

($P\ n$)
and $F: \bigwedge n. n < \text{length } cl \implies \text{inj-on } (F\ n)\ (P\ n)$
shows $\text{inj-on } (\text{UNlift01 } cl\ dl\ P\ F)\ (\text{UNpart1 } cl\ dl\ P)$ (**is** $\text{inj-on } ?G\ ?Q$)
 $\langle \text{proof} \rangle$

lemma *inj-on-singl*:
assumes $\text{inj-on } f\ A$ **and** $a0 \notin A$ **and** $\bigwedge a. a \in A \implies f\ a \neq f\ a0$
shows $\text{inj-on } f\ (\{a0\}\ Un\ A)$
 $\langle \text{proof} \rangle$

lemma *UNlift01-inj-on*:
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\}\ (P\ n) \wedge \{\}\notin P\ n$
and $FP: \bigwedge n. n < \text{length } dl \implies \text{part } \{.. < \text{brn } (dl!n)\}\ (F\ n\ ' (P\ n)) \wedge \{\}\notin F\ n\ ' (P\ n)$
and $F: \bigwedge n. n < \text{length } cl \implies \text{inj-on } (F\ n)\ (P\ n)$
shows $\text{inj-on } (\text{UNlift01 } cl\ dl\ P\ F)\ (\text{UNpart01 } cl\ dl\ P)$
 $\langle \text{proof} \rangle$

lemma *UNlift01-UNpart1*:
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\}\ (P\ n) \wedge \{\}\notin P\ n$
shows $(\text{UNlift01 } cl\ dl\ P\ F)\ ' (\text{UNpart1 } cl\ dl\ P) = \text{UNpart1 } dl\ cl\ (\%n. F\ n\ ' (P\ n))$ (**is** $?G\ ' ?Q = ?R$)
 $\langle \text{proof} \rangle$

lemma *UNlift01-UNpart01*:
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\}\ (P\ n) \wedge \{\}\notin P\ n$
shows $(\text{UNlift01 } cl\ dl\ P\ F)\ ' (\text{UNpart01 } cl\ dl\ P) = \text{UNpart01 } dl\ cl\ (\%n. F\ n\ ' (P\ n))$
 $\langle \text{proof} \rangle$

lemma *emp-UNlift01-UNpart1*:
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\}\ (P\ n) \wedge \{\}\notin P\ n$
and $FP: \bigwedge n. n < \text{length } dl \implies \{\}\notin F\ n\ ' (P\ n)$
shows $\{\}\notin (\text{UNlift01 } cl\ dl\ P\ F)\ ' (\text{UNpart1 } cl\ dl\ P)$ (**is** $\{\}\notin ?R$)
 $\langle \text{proof} \rangle$

lemma *emp-UNlift01-UNpart01*:
assumes $l: \text{length } cl = \text{length } dl$
and $P: \bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\}\ (P\ n) \wedge \{\}\notin P\ n$
and $FP: \bigwedge n. n < \text{length } dl \implies \{\}\notin F\ n\ ' (P\ n)$
shows $\{\}\notin (\text{UNlift01 } cl\ dl\ P\ F)\ ' (\text{UNpart01 } cl\ dl\ P - \{\text{BrnFT } cl\ dl\})$
(is $\{\}\notin ?U\ ' ?V$)
 $\langle \text{proof} \rangle$

lemma *part-UNlift01-UNpart1*:

assumes l : $\text{length } cl = \text{length } dl$ **and** dl : $\text{properL } dl$
and P : $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$
and FP : $\bigwedge n. n < \text{length } dl \implies \text{part } \{..< \text{brn } (dl!n)\} (F \ n \ ' (P \ n))$
shows $\text{part } (\text{BrnNFT } dl \ cl) ((\text{UNlift01 } cl \ dl \ P \ F) \ ' (\text{UNpart1 } cl \ dl \ P))$ (**is** $\text{part } ?C$
 $?R$)
 $\langle \text{proof} \rangle$

lemma $\text{part-UNlift01-UNpart01}$:
assumes l : $\text{length } cl = \text{length } dl$ **and** dl : $\text{properL } dl$
and P : $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$
and FP : $\bigwedge n. n < \text{length } dl \implies \text{part } \{..< \text{brn } (dl!n)\} (F \ n \ ' (P \ n)) \wedge \{\} \notin (F \ n$
 $\ ' (P \ n))$
shows $\text{part } \{..< \text{brnL } dl \ (\text{length } dl)\} ((\text{UNlift01 } cl \ dl \ P \ F) \ ' (\text{UNpart01 } cl \ dl \ P))$
(**is** $\text{part } ?K \ ?R$)
 $\langle \text{proof} \rangle$

lemma diff-frac-eq-1 :
assumes $b \neq (0::\text{real})$
shows $1 - a / b = (b - a) / b$
 $\langle \text{proof} \rangle$

lemma diff-frac-eq-2 :
assumes $b \neq (1::\text{real})$
shows $1 - (a - b) / (1 - b) = (1 - a) / (1 - b)$
(**is** $?L = ?R$)
 $\langle \text{proof} \rangle$

lemma triv-div-mult :
assumes vSF : $vSF \neq (1::\text{real})$
and L : $L = (K - vSF) / (1 - vSF)$ **and** Ln : $L \neq 1$
shows $(VS / (1 - vSF) * V) / (1 - L) = (VS * V) / (1 - K)$
(**is** $?A = ?B$)
 $\langle \text{proof} \rangle$

lemma $\text{ss-wt-ParT-UNlift01}$:
assumes l : $\text{length } cl = \text{length } dl$
and $cldl$: $\text{properL } cl \ \text{properL } dl$ **and** II : $II \in \text{UNpart01 } cl \ dl \ P - \{\text{BrnFT } cl \ dl\}$
and P : $\bigwedge n. n < \text{length } cl \implies \text{part } \{..< \text{brn } (cl!n)\} (P \ n) \wedge \{\} \notin P \ n$
and FP : $\bigwedge n. n < \text{length } dl \implies \text{part } \{..< \text{brn } (dl!n)\} (F \ n \ ' (P \ n))$
and sw :
 $\bigwedge n \ I. \llbracket n < \text{length } cl; I \in P \ n \rrbracket \implies$
 $\quad \text{sum } (wt \ (cl \ ! \ n) \ s) \ I =$
 $\quad \text{sum } (wt \ (dl \ ! \ n) \ t) \ (F \ n \ I)$
and st : $s \approx t$
and $le1$: $\text{sum } (wt \ (\text{ParT } cl) \ s) \ (\text{BrnFT } cl \ dl) < 1$
 $\text{sum } (wt \ (\text{ParT } dl) \ t) \ (\text{BrnFT } dl \ cl) < 1$
shows

$$\begin{aligned}
& \text{sum } (wt \ (ParT \ cl) \ s) \ II \ / \\
& (1 - \text{sum } (wt \ (ParT \ cl) \ s) \ (BrnFT \ cl \ dl)) = \\
& \text{sum } (wt \ (ParT \ dl) \ t) \ (UNlift01 \ cl \ dl \ P \ F \ II) \ / \\
& (1 - \text{sum } (wt \ (ParT \ dl) \ t) \ (BrnFT \ dl \ cl)) \\
& (\text{is } \text{sum } ?vP \ II \ / \ (1 - \text{sum } ?vP \ ?II-0) = \\
& \quad \text{sum } ?wP \ ?JJ \ / \ (1 - \text{sum } ?wP \ ?JJ-0)) \\
& \langle \text{proof} \rangle
\end{aligned}$$

definition *thetaZOParT* where

$$\begin{aligned}
& \text{thetaZOParT} \equiv \\
& \{(ParT \ cl, \ ParT \ dl) \mid \\
& \quad cl \ dl. \\
& \quad \text{properL} \ cl \wedge \text{properL} \ dl \wedge \text{SbisL} \ cl \ dl\}
\end{aligned}$$

lemma *cont-eff-ParT-BrnFT-L*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *cldl*: $\text{properL } cl \ \text{properL } dl \ \text{SbisL } cl \ dl$

and *ii*: $ii \in BrnFT \ cl \ dl$

and *eff-cont*:

$$\begin{aligned}
& \bigwedge n \ I \ i \ j. \llbracket n < \text{length } cl; I \in P \ n; i \in I; j \in F \ n \ I \rrbracket \implies \\
& \quad \text{eff } (cl!n) \ s \ i \approx \text{eff } (dl!n) \ t \ j \wedge \\
& \quad \text{cont } (cl!n) \ s \ i \approx_s \text{cont } (dl!n) \ t \ j
\end{aligned}$$

shows

$$\begin{aligned}
& s \approx \text{eff } (ParT \ cl) \ s \ ii \wedge \\
& \quad (\text{cont } (ParT \ cl) \ s \ ii, \ ParT \ dl) \in \text{thetaZOParT} \\
& (\text{is } ?\text{eff} \wedge ?\text{cont}) \\
& \langle \text{proof} \rangle
\end{aligned}$$

lemma *cont-eff-ParT-BrnFT-R*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *cldl*: $\text{properL } cl \ \text{properL } dl \ \text{SbisL } cl \ dl$

and *jj*: $jj \in BrnFT \ dl \ cl$

and *eff-cont*:

$$\begin{aligned}
& \bigwedge n \ I \ i \ j. \llbracket n < \text{length } cl; I \in P \ n; i \in I; j \in F \ n \ I \rrbracket \implies \\
& \quad \text{eff } (cl!n) \ s \ i \approx \text{eff } (dl!n) \ t \ j \wedge \text{cont } (cl!n) \ s \ i \approx_s \text{cont } (dl!n) \ t \ j
\end{aligned}$$

shows

$$\begin{aligned}
& t \approx \text{eff } (ParT \ dl) \ t \ jj \wedge \\
& \quad (ParT \ cl, \ \text{cont } (ParT \ dl) \ t \ jj) \in \text{thetaZOParT} \\
& (\text{is } ?\text{eff} \wedge ?\text{cont}) \\
& \langle \text{proof} \rangle
\end{aligned}$$

lemma *cont-eff-ParT-UNlift01*:

assumes *l*: $\text{length } cl = \text{length } dl$

and *cldl*: $\text{properL } cl \ \text{properL } dl \ \text{SbisL } cl \ dl$

and *II*: $II \in UNpart01 \ cl \ dl \ P - \{BrnFT \ cl \ dl\}$

and *ii*: $ii \in II$ **and** *jj*: $jj \in UNlift01 \ cl \ dl \ P \ F \ II$

and *P*: $\bigwedge n. n < \text{length } cl \implies \text{part } \{.. < \text{brn } (cl!n)\} \ (P \ n) \wedge \{\} \notin P \ n$

and *FP*: $\bigwedge n. n < \text{length } dl \implies \text{part } \{..< \text{brn } (dl!n)\} (F\ n \text{ ' } (P\ n))$

and *eff-cont*:

$\bigwedge n\ I\ i\ j. \llbracket n < \text{length } cl; I \in P\ n; i \in I; j \in F\ n\ I \rrbracket \implies$

$\text{eff } (cl!n)\ s\ i \approx$

$\text{eff } (dl!n)\ t\ j \wedge$

$\text{cont } (cl!n)\ s\ i \approx s$

$\text{cont } (dl!n)\ t\ j$

and *st*: $s \approx t$

shows

$\text{eff } (ParT\ cl)\ s\ ii \approx \text{eff } (ParT\ dl)\ t\ jj \wedge$

$(\text{cont } (ParT\ cl)\ s\ ii, \text{cont } (ParT\ dl)\ t\ jj) \in \text{thetaZOParT}$

(is $?eff \wedge ?cont$)

<proof>

lemma *thetaZOParT-ZOretr*: $\text{thetaZOParT} \subseteq \text{ZOretr } (\text{thetaZOParT})$

<proof>

lemma *thetaZOParT-ZObis*: $\text{thetaZOParT} \subseteq \text{ZObis}$

<proof>

theorem *ParT-ZObis[simp]*:

assumes *properL cl and properL dl and SbisL cl dl*

shows $ParT\ cl \approx_{01} ParT\ dl$

<proof>

end

end

5 Syntactic Criteria

theory *Syntactic-Criteria*

imports *Compositionality*

begin

context *PL-Indis*

begin

lemma *proper-intros[intro]*:

proper Done

proper (Atm atm)

proper c1 \implies proper c2 \implies proper (Seq c1 c2)

proper c1 \implies proper c2 \implies proper (Ch ch c1 c2)

proper c \implies proper (While tst c)

properL cs \implies proper (Par cs)

properL cs \implies proper (ParT cs)

$(\bigwedge c. c \in \text{set } cs \implies \text{proper } c) \implies cs \neq [] \implies \text{properL } cs$
 ⟨proof⟩

lemma *discr*:

discr Done
presAtm atm \implies discr (Atm atm)
discr c1 \implies discr c2 \implies discr (Seq c1 c2)
discr c1 \implies discr c2 \implies discr (Ch ch c1 c2)
discr c \implies discr (While tst c)
properL cs \implies ($\bigwedge c. c \in \text{set } cs \implies \text{discr } c$) \implies discr (Par cs)
properL cs \implies ($\bigwedge c. c \in \text{set } cs \implies \text{discr } c$) \implies discr (ParT cs)
 ⟨proof⟩

lemma *siso*:

compatAtm atm \implies siso (Atm atm)
siso c1 \implies siso c2 \implies siso (Seq c1 c2)
compatCh ch \implies siso c1 \implies siso c2 \implies siso (Ch ch c1 c2)
compatTst tst \implies siso c \implies siso (While tst c)
properL cs \implies ($\bigwedge c. c \in \text{set } cs \implies \text{siso } c$) \implies siso (Par cs)
properL cs \implies ($\bigwedge c. c \in \text{set } cs \implies \text{siso } c$) \implies siso (ParT cs)
 ⟨proof⟩

lemma *Sbis*:

compatAtm atm \implies Atm atm \approx_s Atm atm
siso c1 \implies c2 \approx_s c2 \implies Seq c1 c2 \approx_s Seq c1 c2
proper c1 \implies proper c2 \implies c1 \approx_s c1 \implies discr c2 \implies Seq c1 c2 \approx_s Seq c1 c2
compatCh ch \implies c1 \approx_s c1 \implies c2 \approx_s c2 \implies Ch ch c1 c2 \approx_s Ch ch c1 c2
properL cs \implies ($\bigwedge c. c \in \text{set } cs \implies c \approx_s c$) \implies Par cs \approx_s Par cs
 ⟨proof⟩

lemma *ZObis*:

compatAtm atm \implies Atm atm \approx_{01} Atm atm
siso c1 \implies c2 \approx_{01} c2 \implies Seq c1 c2 \approx_{01} Seq c1 c2
proper c1 \implies proper c2 \implies c1 \approx_{01} c1 \implies discr c2 \implies Seq c1 c2 \approx_{01} Seq c1 c2
compatCh ch \implies c1 \approx_{01} c1 \implies c2 \approx_{01} c2 \implies Ch ch c1 c2 \approx_{01} Ch ch c1 c2
properL cs \implies ($\bigwedge c. c \in \text{set } cs \implies c \approx_s c$) \implies ParT cs \approx_{01} ParT cs
 ⟨proof⟩

lemma *discr-imp-Sbis*: *proper c \implies discr c \implies c \approx_s c*
 ⟨proof⟩

lemma *siso-imp-Sbis*: *siso c \implies c \approx_s c*
 ⟨proof⟩

lemma *Sbis-imp-ZObis*: *c \approx_s c \implies c \approx_{01} c*
 ⟨proof⟩

fun *SC-discr* **where**

$SC-discr\ Done \iff True$
 $| SC-discr (Atm\ atm) \iff presAtm\ atm$
 $| SC-discr (Seq\ c1\ c2) \iff SC-discr\ c1 \wedge SC-discr\ c2$
 $| SC-discr (Ch\ ch\ c1\ c2) \iff SC-discr\ c1 \wedge SC-discr\ c2$
 $| SC-discr (While\ tst\ c) \iff SC-discr\ c$
 $| SC-discr (ParT\ cs) \iff (\forall c \in set\ cs. SC-discr\ c)$
 $| SC-discr (Par\ cs) \iff (\forall c \in set\ cs. SC-discr\ c)$

theorem *SC-discr-discr*[*intro*]: $proper\ c \implies SC-discr\ c \implies discr\ c$
<proof>

fun *SC-iso* **where**

$SC-iso\ Done \iff True$
 $| SC-iso (Atm\ atm) \iff compatAtm\ atm$
 $| SC-iso (Seq\ c1\ c2) \iff SC-iso\ c1 \wedge SC-iso\ c2$
 $| SC-iso (Ch\ ch\ c1\ c2) \iff compatCh\ ch \wedge SC-iso\ c1 \wedge SC-iso\ c2$
 $| SC-iso (While\ tst\ c) \iff compatTst\ tst \wedge SC-iso\ c$
 $| SC-iso (Par\ cs) \iff (\forall c \in set\ cs. SC-iso\ c)$
 $| SC-iso (ParT\ cs) \iff (\forall c \in set\ cs. SC-iso\ c)$

theorem *SC-iso-iso*[*intro*]: $proper\ c \implies SC-iso\ c \implies iso\ c$
<proof>

fun *SC-Sbis* **where**

$SC-Sbis\ Done \iff True$
 $| SC-Sbis (Atm\ atm) \iff compatAtm\ atm$
 $| SC-Sbis (Seq\ c1\ c2) \iff (SC-iso\ c1 \wedge SC-Sbis\ c2) \vee$
 $(SC-Sbis\ c1 \wedge SC-discr\ c2) \vee$
 $SC-discr (Seq\ c1\ c2) \vee SC-iso (Seq\ c1\ c2)$
 $| SC-Sbis (Ch\ ch\ c1\ c2) \iff (if\ compatCh\ ch$
 $then\ SC-Sbis\ c1 \wedge SC-Sbis\ c2$
 $else\ (SC-discr (Ch\ ch\ c1\ c2) \vee SC-iso (Ch\ ch\ c1\ c2)))$
 $| SC-Sbis (While\ tst\ c) \iff SC-discr (While\ tst\ c) \vee SC-iso (While\ tst\ c)$
 $| SC-Sbis (Par\ cs) \iff (\forall c \in set\ cs. SC-Sbis\ c)$
 $| SC-Sbis (ParT\ cs) \iff SC-iso (ParT\ cs) \vee SC-discr (ParT\ cs)$

theorem *SC-iso-SCbis*[*intro*]: $SC-iso\ c \implies SC-Sbis\ c$
<proof>

theorem *SC-discr-SCbis*[*intro*]: $SC-discr\ c \implies SC-Sbis\ c$
<proof>

declare *SC-iso.simps*[*simp del*]

declare *SC-discr.simps*[*simp del*]


```

declare level.split[split]

instantiation level :: complete-lattice
begin
  definition top-level: top ≡ Hi
  definition bot-level: bot ≡ Lo
  definition inf-level: inf l1 l2 ≡ if Lo ∈ {l1,l2} then Lo else Hi
  definition sup-level: sup l1 l2 ≡ if Hi ∈ {l1,l2} then Hi else Lo
  definition less-eq-level: less-eq l1 l2 ≡ (l1 = Lo ∨ l2 = Hi)
  definition less-level: less l1 l2 ≡ l1 = Lo ∧ l2 = Hi
  definition Inf-level: Inf L ≡ if Lo ∈ L then Lo else Hi
  definition Sup-level: Sup L ≡ if Hi ∈ L then Hi else Lo
instance
  ⟨proof⟩
end

lemma sup-eq-Lo[simp]: sup a b = Lo ↔ a = Lo ∧ b = Lo
  ⟨proof⟩

datatype var = h | h' | l | l'
datatype exp = Ct nat | Var var | Plus exp exp | Minus exp exp
datatype test = Tr | Eq exp exp | Gt exp exp | Non test
datatype atom = Assign var exp
type-synonym choice = real + test
type-synonym state = var ⇒ nat

syntax
  -assign :: 'a ⇒ 'a ⇒ 'a (← ::= → [1000, 61] 61)

syntax-consts
  -assign == Assign

translations
  x ::= expr == CONST Atm (CONST Assign x expr)

primrec sec where
  sec h = Hi
  | sec h' = Hi
  | sec l = Lo
  | sec l' = Lo

fun eval where
  eval (Ct n) s = n
  | eval (Var x) s = s x
  | eval (Plus e1 e2) s = eval e1 s + eval e2 s
  | eval (Minus e1 e2) s = eval e1 s - eval e2 s

fun tval where
  tval Tr s = True

```

$| \text{tval } (Eq \ e1 \ e2) \ s = (\text{eval } e1 \ s = \text{eval } e2 \ s)$
 $| \text{tval } (Gt \ e1 \ e2) \ s = (\text{eval } e1 \ s > \text{eval } e2 \ s)$
 $| \text{tval } (Non \ e) \ s = (\neg \ \text{tval } e \ s)$

fun *aval* **where**

aval (*Assign* *x* *e*) *s* = (*s* (*x* := *eval* *e* *s*))

fun *cval* **where**

$\text{cval } (Inl \ p) \ s = \min \ 1 \ (\max \ 0 \ p)$
 $| \text{cval } (Inr \ \text{tst}) \ s = (\text{if } \text{tval } \text{tst} \ s \ \text{then } 1 \ \text{else } 0)$

definition *indis* :: (*state* * *state*) *setwhere*

indis $\equiv \{(s,t). \ \text{ALL } x. \ \text{sec } x = Lo \ \longrightarrow \ s \ x = t \ x\}$

interpretation *Example-PL: PL-Indis* *aval* *tval* *cval* *indis*

<proof>

fun *exprSec* **where**

$\text{exprSec } (Ct \ n) = Lo$
 $| \text{exprSec } (Var \ x) = \text{sec } x$
 $| \text{exprSec } (Plus \ e1 \ e2) = \sup (\text{exprSec } e1) (\text{exprSec } e2)$
 $| \text{exprSec } (Minus \ e1 \ e2) = \sup (\text{exprSec } e1) (\text{exprSec } e2)$

fun *tstSec* **where**

$\text{tstSec } Tr = Lo$
 $| \text{tstSec } (Eq \ e1 \ e2) = \sup (\text{exprSec } e1) (\text{exprSec } e2)$
 $| \text{tstSec } (Gt \ e1 \ e2) = \sup (\text{exprSec } e1) (\text{exprSec } e2)$
 $| \text{tstSec } (Non \ e) = \text{tstSec } e$

lemma *exprSec-Lo-eval-eq*: $\text{exprSec } \text{expr} = Lo \ \Longrightarrow \ (s, t) \in \text{indis} \ \Longrightarrow \ \text{eval } \text{expr} \ s = \text{eval } \text{expr} \ t$

<proof>

lemma *compatAtmSyntactic[simp]*: $\text{exprSec } \text{expr} = Lo \ \vee \ \text{sec } v = Hi \ \Longrightarrow \ \text{Example-PL.compatAtm } (\text{Assign } v \ \text{expr})$

<proof>

lemma *presAtmSyntactic[simp]*: $\text{sec } v = Hi \ \Longrightarrow \ \text{Example-PL.presAtm } (\text{Assign } v \ \text{expr})$

<proof>

lemma *compatTstSyntactic[simp]*: $\text{tstSec } \text{tst} = Lo \ \Longrightarrow \ \text{Example-PL.compatTst } \text{tst}$

<proof>

lemma *compatPrchSyntactic[simp]*: $\text{Example-PL.compatCh } (Inl \ p)$

<proof>

lemma *compatIfchSyntactic[simp]*: *Example-PL.compatCh (Inr tst) \longleftrightarrow Example-PL.compatTst tst*
 ⟨proof⟩

abbreviation *Ch-half* ($\langle Ch_{1/2} \rangle$) **where** $Ch_{1/2} \equiv Ch (Inl (1/2))$
abbreviation *If* **where** $If\ tst \equiv Ch (Inr\ tst)$

abbreviation *siso* $c \equiv Example-PL.siso\ c$
abbreviation *discr* $c \equiv Example-PL.discr\ c$
abbreviation *Sbis-abbrev* (**infix** $\langle \approx_s \rangle$ 55) **where** $c1 \approx_s c2 \equiv (c1, c2) \in Example-PL.Sbis$
abbreviation *ZObis-abbrev* (**infix** $\langle \approx_{01} \rangle$ 55) **where** $c1 \approx_{01} c2 \equiv (c1, c2) \in Example-PL.ZObis$

abbreviation *SC-siso* $c \equiv Example-PL.SC-siso\ c$
abbreviation *SC-discr* $c \equiv Example-PL.SC-discr\ c$
abbreviation *SC-Sbis* $c \equiv Example-PL.SC-Sbis\ c$
abbreviation *SC-ZObis* $c \equiv Example-PL.SC-ZObis\ c$

lemma *SC-discr (h ::= Ct 0)*
 ⟨proof⟩

6.1 The secure programs from the paper's Example 3

definition [*simp*]: $d0 =$
 $h' ::= Ct\ 0\ ;;$
 $While\ (Gt\ (Var\ h)\ (Ct\ 0))$
 $(Ch_{1/2}\ (h\ ::= Ct\ 0)$
 $(h' ::= Plus\ (Var\ h')\ (Ct\ 1)))$

definition [*simp*]: $d1 =$
 $While\ (Gt\ (Var\ h)\ (Ct\ 0))$
 $(Ch_{1/2}\ (h\ ::= Minus\ (Var\ h)\ (Ct\ 1))$
 $(h\ ::= Plus\ (Var\ h)\ (Ct\ 1)))$

definition [*simp*]: $d2 =$
 $If\ (Eq\ (Var\ l)\ (Ct\ 0))$
 $(l' ::= Ct\ 1)$
 $d0$

definition [*simp*]: $d3 =$
 $h\ ::= Ct\ 5\ ;;$
 $ParT\ [d0,\ (l\ ::= Ct\ 1)]$

theorem *SC-discr d0*
SC-discr d1

SC-Sbis d2
SC-ZObis d2
<proof>

theorem *discr d0*
discr d1
d2 \approx_s d2
d3 \approx_{01} d3
<proof>

end

References

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