

Poincaré Disc Model

Danijela Simić

Filip Marić

Pierre Boutry

April 9, 2025

Abstract

We describe formalization of the Poincaré disc model of hyperbolic geometry within the Isabelle/HOL proof assistant. The model is defined within the extended complex plane (one dimensional complex projective space $\mathbb{C}P^1$), formalized in the AFP entry “Complex Geometry” [6]. Points, lines, congruence of pairs of points, betweenness of triples of points, circles, and isometries are defined within the model. It is shown that the model satisfies all Tarski’s axioms except the Euclid’s axiom. It is shown that it satisfies its negation and the limiting parallels axiom (which proves it to be a model of hyperbolic geometry).

Contents

1	Introduction	3
2	Background theories	3
2.1	Hyperbolic Functions	3
3	Tarski axioms	4
4	H-lines in the Poincaré model	5
4.1	Definition and basic properties of h-lines	5
4.1.1	Collinear points	7
4.1.2	H-lines and inversion	7
4.1.3	Classification of h-lines into Euclidean segments and circles	7
4.1.4	Points on h-line	8
4.1.5	H-line uniqueness	11
4.1.6	H-isometries preserve h-lines	14
4.1.7	Mapping h-lines to x-axis	16
4.2	Construction of the h-line between the two given points	17
4.2.1	Correctness of the construction	20
4.2.2	Existence of h-lines	28
4.2.3	Uniqueness of h-lines	29
4.2.4	Some consequences of line uniqueness	30
4.2.5	Transformations of constructed lines	31
4.2.6	Collinear points and h-lines	32
4.2.7	Points collinear with θ_h	32
4.3	Ideal points of h-lines	34
4.3.1	Calculation of ideal points	34
4.3.2	Ideal points	39
5	H-distance in the Poincaré model	44
5.1	Distance explicit formula	48
5.2	Existence and uniqueness of points with a given distance	54
5.3	Triangle inequality	65
6	H-circles in the Poincaré model	67
6.1	Intersection of circles in special positions	71
6.2	Congruent triangles	74

7	H-betweenness in the Poincaré model	77
7.1	H-betweenness expressed by a cross-ratio	77
7.1.1	H-betweenness is preserved by h-isometries	77
7.1.2	Some elementary properties of h-betweenness	78
7.1.3	H-betweenness and h-collinearity	78
7.1.4	H-betweenness on Euclidean segments	79
7.1.5	H-betweenness and h-collinearity	86
7.2	Some properties of betweenness	87
7.3	Poincare between - sum distances	89
7.4	Some more properties of h-betweenness.	94
8	Intersection of h-lines with the x-axis in the Poincaré model	96
8.1	Betweenness of x-axis intersection	96
8.2	Check if an h-line intersects the x-axis	97
8.3	Check if a Poincaré line intersects the y-axis	100
8.4	Intersection point of a Poincaré line with the x-axis in the unit disc	101
8.5	Check if an h-line intersects the positive part of the x-axis	105
8.6	Check if an h-line intersects the positive part of the y-axis	108
8.7	Position of the intersection point in the unit disc	110
8.8	Ideal points and x-axis intersection	112
9	H-perpendicular h-lines in the Poincaré model	114
10	Poincaré disc model types	124
10.1	H-points	124
10.2	H-lines	124
10.3	H-collinearity	125
10.4	H-isometries	125
10.5	H-distance and h-congruence	126
10.6	H-betweenness	127
11	Poincaré model satisfies Tarski axioms	127
11.1	Pasch axiom	127
11.2	Segment construction axiom	138
11.3	Five segment axiom	140
11.4	Upper dimension axiom	141
11.5	Lower dimension axiom	143
11.6	Negated Euclidean axiom	143
11.7	Continuity axiom	150
11.8	Limiting parallels axiom	157
11.9	Interpretation of locales	170

1 Introduction

Poincaré disc is a model of hyperbolic geometry. That fact has been a mathematical folklore for more than 100 years. However, up to the best of our knowledge, fully precise, formal proofs of this fact are lacking. In this paper we present a formalization of the Poincaré disc model in Isabelle/HOL, introduce its basic notions (h-points, h-lines, h-congruence, h-isometries, h-betweenness) and prove that it models Tarski's axioms except for Euclid's axiom. We show that it satisfies the negation of Euclid's axiom, and, moreover, the existence of limiting parallels axiom. The model is defined within the extended complex plane, which has been described quite precisely by Schwerdfeger [8] and formalized in the previous work of the first two authors [5].

Related work. In 1840 Lobachevsky [3] published developments about non-Euclidean geometry. Hyperbolic geometry is studied through many of its models. The concept of a projective disc model was introduced by Klein while Poincaré investigated the half-plane model proposed by Liouville and Beltrami and primarily studied the isometries of the hyperbolic plane that preserve orientation. In this paper, we focus on the formalization of the latter.

Regarding non-Euclidean geometry, Makarios showed the independence of Euclid's axiom [4]. He did so by formalizing that the Klein–Beltrami model is a model of Tarski's axioms at the exception of Euclid's axiom. Later Coghetto formalized the Klein–Beltrami model within Mizar [1, 2].

2 Background theories

2.1 Hyperbolic Functions

In this section hyperbolic cosine and hyperbolic sine functions are introduced and some of their properties needed for further development are proved.

theory *Hyperbolic-Functions*

imports *Complex-Main Complex-Geometry.More-Complex*
begin

lemma *arcosh-eq-iff*:

fixes $x\ y::\text{real}$
assumes $x \geq 1\ y \geq 1$
shows $\text{arcosh } x = \text{arcosh } y \iff x = y$
by (*smt (verit, best) arcosh-less-iff-real arcosh-real-strict-mono assms*)

lemma *cosh-gt-1 [simp]*:

fixes $x::\text{real}$
assumes $x > 0$
shows $\cosh x > 1$
using *assms cosh-real-strict-mono* **by force**

lemma *cosh-eq-iff*:

fixes $x\ y::\text{real}$
assumes $x \geq 0\ y \geq 0$
shows $\cosh x = \cosh y \iff x = y$
by (*simp add: assms*)

lemma *arcosh-mono*:

fixes $x\ y::\text{real}$
assumes $x \geq 1\ y \geq 1$
shows $\text{arcosh } x \geq \text{arcosh } y \iff x \geq y$
using *arcosh-less-iff-real assms linorder-not-le* **by blast**

lemma *arcosh-add*:

fixes $x\ y::\text{real}$
assumes $x \geq 1\ y \geq 1$
shows $\text{arcosh } x + \text{arcosh } y = \text{arcosh } (x*y + \text{sqrt}((x^2 - 1)*(y^2 - 1)))$
proof–
have $\text{sqrt}((x^2 - 1) * (y^2 - 1)) \geq 0$

```

    using assms
    by simp
  moreover
  have  $x * y \geq 1$ 
    using assms
    by (smt mult-le-cancel-left1)
  ultimately
  have **:  $x * y + \text{sqrt}((x^2 - 1) * (y^2 - 1)) \geq 1$ 
    by linarith
  hence 1:  $0 \leq (x * y + \text{sqrt}((x^2 - 1) * (y^2 - 1)))^2 - 1$ 
    by simp

  have 2:  $x * \text{sqrt}(y^2 - 1) + y * \text{sqrt}(x^2 - 1) \geq 0$ 
    using assms
    by simp

  have  $(x * \text{sqrt}(y^2 - 1) + y * \text{sqrt}(x^2 - 1))^2 = (\text{sqrt}((x * y + \text{sqrt}((x^2 - 1) * (y^2 - 1))))^2 - 1)^2$ 
    using assms
  proof (subst real-sqrt-pow2)
    show  $0 \leq (x * y + \text{sqrt}((x^2 - 1) * (y^2 - 1)))^2 - 1$ 
      by fact
  next
    have  $(x * \text{sqrt}(y^2 - 1))^2 = x^2 * (y^2 - 1)$ 
      by (simp add: <y ≥ 1> power-mult-distrib)
    moreover
    have  $(y * \text{sqrt}(x^2 - 1))^2 = y^2 * (x^2 - 1)$ 
      using assms by (simp add: power-mult-distrib)
    ultimately show  $(x * \text{sqrt}(y^2 - 1) + y * \text{sqrt}(x^2 - 1))^2 = (x * y + \text{sqrt}((x^2 - 1) * (y^2 - 1)))^2 - 1$ 
      using assms
      unfolding power2-sum
      apply (simp add: real-sqrt-mult power-mult-distrib)
      apply (simp add: field-simps)
      done
  qed
  hence  $\text{sqrt}((x * y + \text{sqrt}((x^2 - 1) * (y^2 - 1)))^2 - 1) = x * \text{sqrt}(y^2 - 1) + y * \text{sqrt}(x^2 - 1)$ 
    using power2-eq-iff-nonneg[OF 2 real-sqrt-ge-zero[OF 1]]
    by simp
  thus ?thesis
    using assms **
    apply (simp add: arcosh-real-def)
    apply (subst ln-mult-pos[symmetric])
    apply (smt one-le-power real-sqrt-ge-zero)
    apply (smt (verit) one-le-power real-sqrt-ge-zero)
    by (smt (verit) distrib-right mult.commute real-sqrt-mult)
  qed

```

```

lemma arcosh-double:
  fixes  $x :: \text{real}$ 
  assumes  $x \geq 1$ 
  shows  $2 * \text{arcosh } x = \text{arcosh}(2 * x^2 - 1)$ 
  by (smt arcosh-add arcosh-mono assms one-power2 power2-eq-square real-sqrt-abs)

```

end

3 Tarski axioms

In this section we introduce axioms of Tarski [7] through a series of locales.

```

theory Tarski
imports Main
begin

```

The first locale assumes all Tarski axioms except for the Euclid's axiom and the continuity axiom and corresponds to absolute geometry.

```

locale TarskiAbsolute =

```

```

fixes cong :: 'p ⇒ 'p ⇒ 'p ⇒ 'p ⇒ bool
fixes betw :: 'p ⇒ 'p ⇒ 'p ⇒ bool
assumes cong-reflexive: cong x y y x
assumes cong-transitive: cong x y z u ∧ cong x y v w → cong z u v w
assumes cong-identity: cong x y z z → x = y
assumes segment-construction: ∃ z. betw x y z ∧ cong y z a b
assumes five-segment: x ≠ y ∧ betw x y z ∧ betw x' y' z' ∧ cong x y x' y' ∧ cong y z y' z' ∧ cong x u x' u' ∧ cong
y u y' u' → cong z u z' u'
assumes betw-identity: betw x y x → x = y
assumes Pasch: betw x u z ∧ betw y v z → (∃ a. betw u a y ∧ betw x a v)
assumes lower-dimension: ∃ a. ∃ b. ∃ c. ¬ betw a b c ∧ ¬ betw b c a ∧ ¬ betw c a b
assumes upper-dimension: cong x u x v ∧ cong y u y v ∧ cong z u z v ∧ u ≠ v → betw x y z ∨ betw y z x ∨ betw z
x y
begin

```

The following definitions are used to specify axioms in the following locales.

Point p is on line ab .

definition *on-line where*

on-line $p a b \longleftrightarrow \text{betw } p a b \vee \text{betw } a p b \vee \text{betw } a b p$

Point p is on ray ab .

definition *on-ray where*

on-ray $p a b \longleftrightarrow \text{betw } a p b \vee \text{betw } a b p$

Point p is inside angle abc .

definition *in-angle where*

in-angle $p a b c \longleftrightarrow b \neq a \wedge b \neq c \wedge p \neq b \wedge (\exists x. \text{betw } a x c \wedge x \neq a \wedge x \neq c \wedge \text{on-ray } p b x)$

Ray $r_a r_b$ meets the line $l_a l_b$.

definition *ray-meets-line where*

ray-meets-line $ra rb la lb \longleftrightarrow (\exists x. \text{on-ray } x ra rb \wedge \text{on-line } x la lb)$

end

The second locale adds the negation of Euclid's axiom and limiting parallels and corresponds to hyperbolic geometry.

locale *TarskiHyperbolic = TarskiAbsolute +*

assumes *euclid-negation*: ∃ $a b c d t$. $\text{betw } a d t \wedge \text{betw } b d c \wedge a \neq d \wedge (\forall x y. \text{betw } a b x \wedge \text{betw } a c y \rightarrow \neg \text{betw } x t y)$

assumes *limiting-parallels*: ¬ *on-line* $a x_1 x_2 \implies$

(∃ $a_1 a_2$. ¬ *on-line* $a a_1 a_2 \wedge$

¬ *ray-meets-line* $a a_1 x_1 x_2 \wedge$

¬ *ray-meets-line* $a a_2 x_1 x_2 \wedge$

(∀ a' . *in-angle* $a' a_1 a a_2 \rightarrow \text{ray-meets-line } a a' x_1 x_2$))

The third locale adds the continuity axiom and corresponds to elementary hyperbolic geometry.

locale *ElementaryTarskiHyperbolic = TarskiHyperbolic +*

assumes *continuity*: $\llbracket \exists a. \forall x. \forall y. \varphi x \wedge \psi y \rightarrow \text{betw } a x y \rrbracket \implies \exists b. \forall x. \forall y. \varphi x \wedge \psi y \rightarrow \text{betw } x b y$

end

4 H-lines in the Poincaré model

theory *Poincare-Lines*

imports *Complex-Geometry.Unit-Circle-Preserving-Moebius Complex-Geometry.Circlines-Angle*

begin

4.1 Definition and basic properties of h-lines

H-lines in the Poincaré model are either line segments passing through the origin or segments (within the unit disc) of circles that are perpendicular to the unit circle. Algebraically these are circlines that are represented by Hermitean matrices of the form

$$H = \begin{pmatrix} A & B \\ \overline{B} & A \end{pmatrix},$$

for $A \in \mathbb{R}$, and $B \in \mathbb{C}$, and $|B|^2 > A^2$, where the circline equation is the usual one: $z^*Hz = 0$, for homogenous coordinates z .

definition *is-poincare-line-cmat* :: *complex-mat* \Rightarrow *bool* **where**
 [simp]: *is-poincare-line-cmat* $H \iff$
 (let $(A, B, C, D) = H$
 in *hermitean* $(A, B, C, D) \wedge A = D \wedge (\text{cmod } B)^2 > (\text{cmod } A)^2$)

lift-definition *is-poincare-line-clmat* :: *circline-mat* \Rightarrow *bool* **is** *is-poincare-line-cmat*
done

We introduce the predicate that checks if a given complex matrix is a matrix of a h-line in the Poincaré model, and then by means of the lifting package lift it to the type of non-zero Hermitean matrices, and then to circlines (that are equivalence classes of such matrices).

lift-definition *is-poincare-line* :: *circline* \Rightarrow *bool* **is** *is-poincare-line-clmat*

proof (*transfer*, *transfer*)
fix $H1\ H2$:: *complex-mat*
assume hh : *hermitean* $H1 \wedge H1 \neq \text{mat-zero}$ *hermitean* $H2 \wedge H2 \neq \text{mat-zero}$
assume *circline-eq-cmat* $H1\ H2$
thus *is-poincare-line-cmat* $H1 \iff$ *is-poincare-line-cmat* $H2$
using hh
by (*cases* $H1$, *cases* $H2$) (*auto simp: norm-mult power-mult-distrib*)
qed

lemma *is-poincare-line-mk-circline*:
assumes $(A, B, C, D) \in \text{hermitean-nonzero}$
shows *is-poincare-line* $(\text{mk-circline } A\ B\ C\ D) \iff (\text{cmod } B)^2 > (\text{cmod } A)^2 \wedge A = D$
using *assms*
by (*transfer*, *transfer*, *auto simp add: Let-def*)

Abstract characterisation of *is-poincare-line* predicate: H-lines in the Poincaré model are real circlines (circlines with the negative determinant) perpendicular to the unit circle.

lemma *is-poincare-line-iff*:
shows *is-poincare-line* $H \iff$ *circline-type* $H = -1 \wedge$ *perpendicular* H *unit-circle*
unfolding *perpendicular-def*
proof (*simp*, *transfer*, *transfer*)
fix H
assume hh : *hermitean* $H \wedge H \neq \text{mat-zero}$
obtain $A\ B\ C\ D$ **where** $*$: $H = (A, B, C, D)$
by (*cases* H , *auto*)
have $**$: *is-real* A *is-real* $D\ C = \text{cnj } B$
using $hh\ *$ *hermitean-elems*
by *auto*
hence $(\text{Re } A = \text{Re } D \wedge \text{cmod } A * \text{cmod } A < \text{cmod } B * \text{cmod } B) =$
 $(\text{Re } A * \text{Re } D < \text{Re } B * \text{Re } B + \text{Im } B * \text{Im } B \wedge (\text{Re } D = \text{Re } A \vee \text{Re } A * \text{Re } D = \text{Re } B * \text{Re } B + \text{Im } B * \text{Im } B))$
using $*$
by (*smt cmod-power2 power2-eq-square zero-power2*)
thus *is-poincare-line-cmat* $H \iff$
circline-type-cmat $H = -1 \wedge$ *cos-angle-cmat* (*of-circline-cmat* H) *unit-circle-cmat* $= 0$
using $**$
by (*auto simp add: sgn-1-neg complex-eq-if-Re-eq cmod-square power2-eq-square simp del: pos-oriented-cmat-def*)
qed

The *x-axis* is an h-line.

lemma *is-poincare-line-x-axis* [*simp*]:
shows *is-poincare-line* *x-axis*
by (*transfer*, *transfer*) (*auto simp add: hermitean-def mat-adj-def mat-cnj-def*)

The *unit-circle* is not an h-line.

lemma *not-is-poincare-line-unit-circle* [*simp*]:
shows \neg *is-poincare-line* *unit-circle*
by (*transfer*, *transfer*, *simp*)

4.1.1 Collinear points

Points are collinear if they all belong to an h-line.

definition *poincare-collinear* :: complex-homo set \Rightarrow bool **where**
poincare-collinear $S \iff (\exists p. \text{is-poincare-line } p \wedge S \subseteq \text{circline-set } p)$

4.1.2 H-lines and inversion

Every h-line in the Poincaré model contains the inverse (wrt. the unit circle) of each of its points (note that at most one of them belongs to the unit disc).

lemma *is-poincare-line-inverse-point*:

assumes *is-poincare-line* H $u \in \text{circline-set } H$

shows *inversion* $u \in \text{circline-set } H$

using *assms*

unfolding *is-poincare-line-iff* *circline-set-def* *perpendicular-def* *inversion-def*

apply *simp*

proof (*transfer*, *transfer*)

fix u H

assume *hh*: *hermitean* $H \wedge H \neq \text{mat-zero } u \neq \text{vec-zero}$ **and**

aa: *circline-type-cmat* $H = -1 \wedge \text{cos-angle-cmat (of-circline-cmat } H) \text{ unit-circle-cmat} = 0$ *on-circline-cmat-cvec*

H u

obtain A B C D $u1$ $u2$ **where** $*$: $H = (A, B, C, D)$ $u = (u1, u2)$

by (*cases* H , *cases* u , *auto*)

have *is-real* A *is-real* D $C = \text{cnj } B$

using $*$ *hh* *hermitean-elems*

by *auto*

moreover

have $A = D$

using *aa*(1) $*$ *is-real* A *is-real* D

by (*auto simp del: pos-oriented-cmat-def simp add: complex.expand split: if-split-asm*)

thus *on-circline-cmat-cvec* H (*conjugate-cvec (reciprocal-cvec* $u)$)

using *aa*(2) $*$

by (*simp add: vec-cnj-def field-simps*)

qed

Every h-line in the Poincaré model and is invariant under unit circle inversion.

lemma *circline-inversion-poincare-line*:

assumes *is-poincare-line* H

shows *circline-inversion* $H = H$

proof –

obtain u v w **where** $*$: $u \neq v$ $v \neq w$ $u \neq w$ $\{u, v, w\} \subseteq \text{circline-set } H$

using *assms* *is-poincare-line-iff*[*of* H]

using *circline-type-neg-card-gt3*[*of* H]

by *auto*

hence $\{\text{inversion } u, \text{inversion } v, \text{inversion } w\} \subseteq \text{circline-set (circline-inversion } H)$

$\{\text{inversion } u, \text{inversion } v, \text{inversion } w\} \subseteq \text{circline-set } H$

using *is-poincare-line-inverse-point*[*OF* *assms*]

by *auto*

thus *?thesis*

using $*$ *unique-circline-set*[*of* *inversion* u *inversion* v *inversion* w]

by (*metis insert-subset inversion-involution*)

qed

4.1.3 Classification of h-lines into Euclidean segments and circles

If an h-line contains zero, than it also contains infinity (the inverse point of zero) and is by definition an Euclidean line.

lemma *is-poincare-line-trough-zero-trough-infty* [*simp*]:

assumes *is-poincare-line* l **and** $0_h \in \text{circline-set } l$

shows $\infty_h \in \text{circline-set } l$

using *is-poincare-line-inverse-point*[*OF* *assms*]

by *simp*

lemma *is-poincare-line-trough-zero-is-line*:

```

assumes is-poincare-line l and  $0_h \in \text{circline-set } l$ 
shows is-line l
using assms
using inf-in-circline-set is-poincare-line-trough-zero-trough-infty
by blast

```

If an h-line does not contain zero, than it also does not contain infinity (the inverse point of zero) and is by definition an Euclidean circle.

```

lemma is-poincare-line-not-trough-zero-not-trough-infty [simp]:
assumes is-poincare-line l
assumes  $0_h \notin \text{circline-set } l$ 
shows  $\infty_h \notin \text{circline-set } l$ 
using assms
using is-poincare-line-inverse-point[OF assms(1), of  $\infty_h$ ]
by auto

```

```

lemma is-poincare-line-not-trough-zero-is-circle:
assumes is-poincare-line l  $0_h \notin \text{circline-set } l$ 
shows is-circle l
using assms
using inf-in-circline-set is-poincare-line-not-trough-zero-not-trough-infty
by auto

```

4.1.4 Points on h-line

Each h-line in the Poincaré model contains at least two different points within the unit disc.

First we prove an auxiliary lemma.

```

lemma ex-is-poincare-line-points':
assumes i12:  $i1 \in \text{circline-set } H \cap \text{unit-circle-set}$ 
            $i2 \in \text{circline-set } H \cap \text{unit-circle-set}$ 
            $i1 \neq i2$ 
assumes a:  $a \in \text{circline-set } H$   $a \notin \text{unit-circle-set}$ 
shows  $\exists b. b \neq i1 \wedge b \neq i2 \wedge b \neq a \wedge b \neq \text{inversion } a \wedge b \in \text{circline-set } H$ 
proof –
have inversion a  $\notin \text{unit-circle-set}$ 
using  $\langle a \notin \text{unit-circle-set} \rangle$ 
unfolding unit-circle-set-def circline-set-def
by (metis inversion-id-iff-on-unit-circle inversion-involution mem-Collect-eq)

have  $a \neq \text{inversion } a$ 
using  $\langle a \notin \text{unit-circle-set} \rangle$  inversion-id-iff-on-unit-circle[of a]
unfolding unit-circle-set-def circline-set-def
by auto

have  $a \neq i1$   $a \neq i2$   $\text{inversion } a \neq i1$   $\text{inversion } a \neq i2$ 
using assms  $\langle \text{inversion } a \notin \text{unit-circle-set} \rangle$ 
by auto

then obtain b where cr2:  $\text{cross-ratio } b \ i1 \ a \ i2 = \text{of-complex } 2$ 
using  $\langle i1 \neq i2 \rangle$ 
using ex-cross-ratio[of i1 a i2]
by blast

have distinct-b:  $b \neq i1$   $b \neq i2$   $b \neq a$ 
using  $\langle i1 \neq i2 \rangle \langle a \neq i1 \rangle \langle a \neq i2 \rangle$ 
using ex1-cross-ratio[of i1 a i2]
using cross-ratio-0[of i1 a i2] cross-ratio-1[of i1 a i2] cross-ratio-inf[of i1 i2 a]
using cr2
by auto

hence  $b \in \text{circline-set } H$ 
using assms four-points-on-circline-iff-cross-ratio-real[of b i1 a i2] cr2
using unique-circline-set[of i1 i2 a]
by auto

```


moreover

have $b \neq \text{inversion } a$

proof (rule ccontr)

assume *: $\neg ?thesis$

have $\text{inversion } i1 = i1 \text{ inversion } i2 = i2$

using $i12$

unfolding $\text{unit-circle-set-def}$

by $auto$

hence $\text{cross-ratio } (\text{inversion } a) \ i1 \ a \ i2 = \text{cross-ratio } a \ i1 \ (\text{inversion } a) \ i2$

using * $\text{cross-ratio-inversion}[of \ i1 \ a \ i2 \ b] \ \langle a \neq i1 \rangle \ \langle a \neq i2 \rangle \ \langle i1 \neq i2 \rangle \ \langle b \neq i1 \rangle$

using $\text{four-points-on-circline-iff-cross-ratio-real}[of \ b \ i1 \ a \ i2]$

using $i12 \ \text{distinct-b-conjugate-id-iff}[of \ \text{cross-ratio } b \ i1 \ a \ i2]$

using $i12 \ a \ \langle b \in \text{circline-set } H \rangle$

by $auto$

hence $\text{cross-ratio } (\text{inversion } a) \ i1 \ a \ i2 \neq \text{of-complex } 2$

using $\text{cross-ratio-commute-13}[of \ \text{inversion } a \ i1 \ a \ i2]$

using reciprocal-id-iff

using of-complex-inj

by $force$

thus $False$

using * $cr2$

by simp

qed

ultimately

show $?thesis$

using $\text{assms } \langle b \neq i1 \rangle \ \langle b \neq i2 \rangle \ \langle b \neq a \rangle$

by $auto$

qed

Now we can prove the statement.

lemma $\text{ex-is-poincare-line-points}$:

assumes $\text{is-poincare-line } H$

shows $\exists \ u \ v. \ u \in \text{unit-disc} \wedge v \in \text{unit-disc} \wedge u \neq v \wedge \{u, v\} \subseteq \text{circline-set } H$

proof—

obtain $u \ v \ w$ where *: $u \neq v \ v \neq w \ u \neq w \ \{u, v, w\} \subseteq \text{circline-set } H$

using $\text{assms } \text{is-poincare-line-iff}[of \ H]$

using $\text{circline-type-neg-card-gt3}[of \ H]$

by $auto$

have $\neg \{u, v, w\} \subseteq \text{unit-circle-set}$

using $\text{unique-circline-set}[of \ u \ v \ w] \ *$

by $(\text{metis } \text{assms } \text{insert-subset } \text{not-is-poincare-line-unit-circle } \text{unit-circle-set-def})$

hence $H \neq \text{unit-circle}$

unfolding $\text{unit-circle-set-def}$

using *

by $auto$

show $?thesis$

proof (cases $(u \in \text{unit-disc} \wedge v \in \text{unit-disc}) \vee$

$(u \in \text{unit-disc} \wedge w \in \text{unit-disc}) \vee$

$(v \in \text{unit-disc} \wedge w \in \text{unit-disc}))$)

case $True$

thus $?thesis$

using *

by $auto$

next

case $False$

have $\exists \ a \ b. \ a \neq b \wedge a \neq \text{inversion } b \wedge a \in \text{circline-set } H \wedge b \in \text{circline-set } H \wedge a \notin \text{unit-circle-set} \wedge b \notin \text{unit-circle-set}$

proof (cases $(u \in \text{unit-circle-set} \wedge v \in \text{unit-circle-set}) \vee$

$(u \in \text{unit-circle-set} \wedge w \in \text{unit-circle-set}) \vee$

$(v \in \text{unit-circle-set} \wedge w \in \text{unit-circle-set}))$)

```

case True
then obtain i1 i2 a where *:
  i1 ∈ unit-circle-set ∩ circline-set H i2 ∈ unit-circle-set ∩ circline-set H
  a ∈ circline-set H a ∉ unit-circle-set
  i1 ≠ i2 i1 ≠ a i2 ≠ a
  using * ⟨¬ {u, v, w} ⊆ unit-circle-set⟩
  by auto
then obtain b where b ∈ circline-set H b ≠ i1 b ≠ i2 b ≠ a b ≠ inversion a
  using ex-is-poincare-line-points'[of i1 H i2 a]
  by blast

hence b ∉ unit-circle-set
  using * ⟨H ≠ unit-circle⟩ unique-circline-set[of i1 i2 b]
  unfolding unit-circle-set-def
  by auto

thus ?thesis
  using * ⟨b ∈ circline-set H⟩ ⟨b ≠ a⟩ ⟨b ≠ inversion a⟩
  by auto
next
case False
then obtain f g h where
  *: f ≠ g f ∈ circline-set H f ∉ unit-circle-set
     g ∈ circline-set H g ∉ unit-circle-set
     h ∈ circline-set H h ≠ f h ≠ g
  using *
  by auto
show ?thesis
proof (cases f = inversion g)
case False
  thus ?thesis
    using *
    by auto
next
case True
  show ?thesis
  proof (cases h ∈ unit-circle-set)
  case False
  thus ?thesis
    using * ⟨f = inversion g⟩
    by auto
next
case True
  obtain m where cr2: cross-ratio m h f g = of-complex 2
    using ex-cross-ratio[of h f g] * ⟨f ≠ g⟩ ⟨h ≠ f⟩ ⟨h ≠ g⟩
    by auto
  hence m ≠ h m ≠ f m ≠ g
    using ⟨h ≠ f⟩ ⟨h ≠ g⟩ ⟨f ≠ g⟩
    using ex1-cross-ratio[of h f g]
    using cross-ratio-0[of h f g] cross-ratio-1[of h f g] cross-ratio-inf[of h g f]
    using cr2
    by auto
  hence m ∈ circline-set H
    using four-points-on-circline-iff-cross-ratio-real[of m h f g] cr2
    using ⟨h ≠ f⟩ ⟨h ≠ g⟩ ⟨f ≠ g⟩ *
    using unique-circline-set[of h f g]
    by auto

show ?thesis
proof (cases m ∈ unit-circle-set)
case False
  thus ?thesis
    using ⟨m ≠ f⟩ ⟨m ≠ g⟩ ⟨f = inversion g⟩ * ⟨m ∈ circline-set H⟩
    by auto
next
case True

```

```

then obtain n where n ≠ h n ≠ m n ≠ f n ≠ inversion f n ∈ circline-set H
  using * ex-is-poincare-line-points'[of h H m f] * ⟨m ∈ circline-set H⟩ ⟨h ∈ unit-circle-set⟩ ⟨m ≠ h⟩
  by auto
hence n ∉ unit-circle-set
  using * ⟨H ≠ unit-circle⟩ unique-circline-set[of m n h]
  using ⟨m ≠ h⟩ ⟨m ∈ unit-circle-set⟩ ⟨h ∈ unit-circle-set⟩ ⟨m ∈ circline-set H⟩
  unfolding unit-circle-set-def
  by auto

thus ?thesis
  using * ⟨n ∈ circline-set H⟩ ⟨n ≠ f⟩ ⟨n ≠ inversion f⟩
  by auto
qed
qed
qed
qed
then obtain a b where ab: a ≠ b a ≠ inversion b a ∈ circline-set H b ∈ circline-set H a ∉ unit-circle-set b ∉
unit-circle-set
  by blast
have ∀ x. x ∈ circline-set H ∧ x ∉ unit-circle-set → (∃ x'. x' ∈ circline-set H ∩ unit-disc ∧ (x' = x ∨ x' =
inversion x))
proof safe
  fix x
  assume x: x ∈ circline-set H x ∉ unit-circle-set
  show ∃ x'. x' ∈ circline-set H ∩ unit-disc ∧ (x' = x ∨ x' = inversion x)
  proof (cases x ∈ unit-disc)
    case True
    thus ?thesis
      using x
      by auto
    next
    case False
    hence x ∈ unit-disc-compl
      using x in-on-out-univ[of ounit-circle]
      unfolding unit-circle-set-def unit-disc-def unit-disc-compl-def
      by auto
    hence inversion x ∈ unit-disc
      using inversion-unit-disc-compl
      by blast
    thus ?thesis
      using is-poincare-line-inverse-point[OF assms, of x] x
      by auto
  qed
qed
then obtain a' b' where
  *: a' ∈ circline-set H a' ∈ unit-disc b' ∈ circline-set H b' ∈ unit-disc and
  **: a' = a ∨ a' = inversion a b' = b ∨ b' = inversion b
  using ab
  by blast
have a' ≠ b'
  using ⟨a ≠ b⟩ ⟨a ≠ inversion b⟩ ** *
  by (metis inversion-involution)
thus ?thesis
  using *
  by auto
qed
qed

```

4.1.5 H-line uniqueness

There is no more than one h-line that contains two different h-points (in the disc).

lemma *unique-is-poincare-line*:

```

assumes in-disc: u ∈ unit-disc v ∈ unit-disc u ≠ v
assumes pl: is-poincare-line l1 is-poincare-line l2
assumes on-l: {u, v} ⊆ circline-set l1 ∩ circline-set l2
shows l1 = l2

```

```

proof-
  have  $u \neq \text{inversion } u \ v \neq \text{inversion } u$ 
    using in-disc
    using inversion-noteq-unit-disc[of  $u \ v$ ]
    using inversion-noteq-unit-disc[of  $u \ u$ ]
    by auto
  thus ?thesis
    using on-l
    using unique-circline-set[of  $u \ \text{inversion } u \ v$ ]  $\langle u \neq v \rangle$ 
    using is-poincare-line-inverse-point[of  $l1 \ u$ ]
    using is-poincare-line-inverse-point[of  $l2 \ u$ ]
    using pl
    by auto
qed

```

For the rest of our formalization it is often useful to consider points on h-lines that are not within the unit disc. Many lemmas in the rest of this section will have such generalizations.

There is no more than one h-line that contains two different and not mutually inverse points (not necessary in the unit disc).

lemma *unique-is-poincare-line-general*:

```

assumes different:  $u \neq v \ u \neq \text{inversion } v$ 
assumes pl: is-poincare-line  $l1$  is-poincare-line  $l2$ 
assumes on-l:  $\{u, v\} \subseteq \text{circline-set } l1 \cap \text{circline-set } l2$ 
shows  $l1 = l2$ 

```

```

proof (cases  $u \neq \text{inversion } u$ )
  case True
  thus ?thesis
    using unique-circline-set[of  $u \ \text{inversion } u \ v$ ]
    using assms
    using is-poincare-line-inverse-point by force

```

```

next
  case False
  show ?thesis
  proof (cases  $v \neq \text{inversion } v$ )
    case True
    thus ?thesis
      using unique-circline-set[of  $u \ \text{inversion } v \ v$ ]
      using assms
      using is-poincare-line-inverse-point by force

```

```

next
  case False

```

```

have on-circline unit-circle  $u$  on-circline unit-circle  $v$ 
  using  $\langle \neg u \neq \text{inversion } u \ \neg v \neq \text{inversion } v \rangle$ 
  using inversion-id-iff-on-unit-circle
  by fastforce+

```

```

thus ?thesis
  using pl on-l  $\langle u \neq v \rangle$ 
  unfolding circline-set-def
  apply simp

```

```

proof (transfer, transfer, safe)
  fix  $u1 \ u2 \ v1 \ v2 \ A1 \ B1 \ C1 \ D1 \ A2 \ B2 \ C2 \ D2 :: \text{complex}$ 
  let  $?u = (u1, u2)$  and  $?v = (v1, v2)$  and  $?H1 = (A1, B1, C1, D1)$  and  $?H2 = (A2, B2, C2, D2)$ 
  assume *:  $?u \neq \text{vec-zero} \ ?v \neq \text{vec-zero}$ 

```

```

  on-circline-cmat-cvec unit-circle-cmat  $?u$  on-circline-cmat-cvec unit-circle-cmat  $?v$ 
  is-poincare-line-cmat  $?H1$  is-poincare-line-cmat  $?H2$ 
  hermitean  $?H1$   $?H1 \neq \text{mat-zero}$  hermitean  $?H2$   $?H2 \neq \text{mat-zero}$ 
  on-circline-cmat-cvec  $?H1 \ ?u$  on-circline-cmat-cvec  $?H1 \ ?v$ 
  on-circline-cmat-cvec  $?H2 \ ?u$  on-circline-cmat-cvec  $?H2 \ ?v$ 
   $\neg (u1, u2) \approx_v (v1, v2)$ 

```

```

have **:  $A1 = D1 \ A2 = D2 \ C1 = \text{cnj } B1 \ C2 = \text{cnj } B2$  is-real  $A1$  is-real  $A2$ 
  using  $\langle \text{is-poincare-line-cmat } ?H1 \rangle \langle \text{is-poincare-line-cmat } ?H2 \rangle$ 
  using  $\langle \text{hermitean } ?H1 \rangle \langle ?H1 \neq \text{mat-zero} \rangle \langle \text{hermitean } ?H2 \rangle \langle ?H2 \neq \text{mat-zero} \rangle$ 
  using hermitean-elems
  by auto

```

have $uv: u1 \neq 0 \ u2 \neq 0 \ v1 \neq 0 \ v2 \neq 0$
using $*(1-4)$
by $(auto \ simp \ add: \ vec\text{-}cnj\text{-}def)$

have $u: \text{cor}((\text{Re}(u1/u2))^2) + \text{cor}((\text{Im}(u1/u2))^2) = 1$
using $\langle on\text{-}circline\text{-}cmat\text{-}cvec \ unit\text{-}circle\text{-}cmat \ ?u \rangle \ uv$
apply $(subst \ of\text{-}real\text{-}add[symmetric])$
apply $(subst \ complex\text{-}mult\text{-}cnj[symmetric])$
apply $(simp \ add: \ vec\text{-}cnj\text{-}def \ mult.commute)$
done

have $v: \text{cor}((\text{Re}(v1/v2))^2) + \text{cor}((\text{Im}(v1/v2))^2) = 1$
using $\langle on\text{-}circline\text{-}cmat\text{-}cvec \ unit\text{-}circle\text{-}cmat \ ?v \rangle \ uv$
apply $(subst \ of\text{-}real\text{-}add[symmetric])$
apply $(subst \ complex\text{-}mult\text{-}cnj[symmetric])$
apply $(simp \ add: \ vec\text{-}cnj\text{-}def \ mult.commute)$
done

have

$A1 * (\text{cor}((\text{Re}(u1/u2))^2) + \text{cor}((\text{Im}(u1/u2))^2) + 1) + \text{cor}(\text{Re } B1) * \text{cor}(2 * \text{Re}(u1/u2)) + \text{cor}(\text{Im } B1)$
 $* \text{cor}(2 * \text{Im}(u1/u2)) = 0$
 $A2 * (\text{cor}((\text{Re}(u1/u2))^2) + \text{cor}((\text{Im}(u1/u2))^2) + 1) + \text{cor}(\text{Re } B2) * \text{cor}(2 * \text{Re}(u1/u2)) + \text{cor}(\text{Im } B2)$
 $* \text{cor}(2 * \text{Im}(u1/u2)) = 0$
 $A1 * (\text{cor}((\text{Re}(v1/v2))^2) + \text{cor}((\text{Im}(v1/v2))^2) + 1) + \text{cor}(\text{Re } B1) * \text{cor}(2 * \text{Re}(v1/v2)) + \text{cor}(\text{Im } B1) *$
 $\text{cor}(2 * \text{Im}(v1/v2)) = 0$
 $A2 * (\text{cor}((\text{Re}(v1/v2))^2) + \text{cor}((\text{Im}(v1/v2))^2) + 1) + \text{cor}(\text{Re } B2) * \text{cor}(2 * \text{Re}(v1/v2)) + \text{cor}(\text{Im } B2) *$
 $\text{cor}(2 * \text{Im}(v1/v2)) = 0$
using $\text{circline-equation-quadratic-equation}[of \ A1 \ u1/u2 \ B1 \ D1 \ \text{Re}(u1/u2) \ \text{Im}(u1/u2) \ \text{Re } B1 \ \text{Im } B1]$
using $\text{circline-equation-quadratic-equation}[of \ A2 \ u1/u2 \ B2 \ D2 \ \text{Re}(u1/u2) \ \text{Im}(u1/u2) \ \text{Re } B2 \ \text{Im } B2]$
using $\text{circline-equation-quadratic-equation}[of \ A1 \ v1/v2 \ B1 \ D1 \ \text{Re}(v1/v2) \ \text{Im}(v1/v2) \ \text{Re } B1 \ \text{Im } B1]$
using $\text{circline-equation-quadratic-equation}[of \ A2 \ v1/v2 \ B2 \ D2 \ \text{Re}(v1/v2) \ \text{Im}(v1/v2) \ \text{Re } B2 \ \text{Im } B2]$
using $\langle on\text{-}circline\text{-}cmat\text{-}cvec \ ?H1 \ ?u \rangle \ \langle on\text{-}circline\text{-}cmat\text{-}cvec \ ?H2 \ ?u \rangle$
using $\langle on\text{-}circline\text{-}cmat\text{-}cvec \ ?H1 \ ?v \rangle \ \langle on\text{-}circline\text{-}cmat\text{-}cvec \ ?H2 \ ?v \rangle$
using $** \ uv$
by $(simp\text{-}all \ add: \ vec\text{-}cnj\text{-}def \ field\text{-}simps)$

hence

$A1 + \text{cor}(\text{Re } B1) * \text{cor}(\text{Re}(u1/u2)) + \text{cor}(\text{Im } B1) * \text{cor}(\text{Im}(u1/u2)) = 0$
 $A1 + \text{cor}(\text{Re } B1) * \text{cor}(\text{Re}(v1/v2)) + \text{cor}(\text{Im } B1) * \text{cor}(\text{Im}(v1/v2)) = 0$
 $A2 + \text{cor}(\text{Re } B2) * \text{cor}(\text{Re}(u1/u2)) + \text{cor}(\text{Im } B2) * \text{cor}(\text{Im}(u1/u2)) = 0$
 $A2 + \text{cor}(\text{Re } B2) * \text{cor}(\text{Re}(v1/v2)) + \text{cor}(\text{Im } B2) * \text{cor}(\text{Im}(v1/v2)) = 0$
using $u \ v$
by $simp\text{-}all \ algebra+$

hence

$\text{cor}(\text{Re } A1 + \text{Re } B1 * \text{Re}(u1/u2) + \text{Im } B1 * \text{Im}(u1/u2)) = 0$
 $\text{cor}(\text{Re } A2 + \text{Re } B2 * \text{Re}(u1/u2) + \text{Im } B2 * \text{Im}(u1/u2)) = 0$
 $\text{cor}(\text{Re } A1 + \text{Re } B1 * \text{Re}(v1/v2) + \text{Im } B1 * \text{Im}(v1/v2)) = 0$
 $\text{cor}(\text{Re } A2 + \text{Re } B2 * \text{Re}(v1/v2) + \text{Im } B2 * \text{Im}(v1/v2)) = 0$
using $\langle is\text{-}real \ A1 \rangle \ \langle is\text{-}real \ A2 \rangle$
by $simp\text{-}all$

hence

$\text{Re } A1 + \text{Re } B1 * \text{Re}(u1/u2) + \text{Im } B1 * \text{Im}(u1/u2) = 0$
 $\text{Re } A1 + \text{Re } B1 * \text{Re}(v1/v2) + \text{Im } B1 * \text{Im}(v1/v2) = 0$
 $\text{Re } A2 + \text{Re } B2 * \text{Re}(u1/u2) + \text{Im } B2 * \text{Im}(u1/u2) = 0$
 $\text{Re } A2 + \text{Re } B2 * \text{Re}(v1/v2) + \text{Im } B2 * \text{Im}(v1/v2) = 0$
using $of\text{-}real\text{-}eq\text{-}0\text{-}iff$
by $blast+$

moreover

have $\text{Re}(u1/u2) \neq \text{Re}(v1/v2) \vee \text{Im}(u1/u2) \neq \text{Im}(v1/v2)$
proof $(rule \ ccontr)$
assume $\neg \ ?thesis$

hence $u1/u2 = v1/v2$
using *complex-eqI* **by** *blast*
thus *False*
using $uv \langle \neg (u1, u2) \approx_v (v1, v2) \rangle$
using $*(1) *(2)$ *complex-cvec-eq-mix*[*OF* $*(1) *(2)$]
by (*auto simp add: field-simps*)
qed

moreover

have $Re\ A1 \neq 0 \vee Re\ B1 \neq 0 \vee Im\ B1 \neq 0$
using $\langle ?H1 \neq mat-zero \rangle **$
by (*metis complex-cnj-zero complex-of-real-Re mat-zero-def of-real-0*)

ultimately

obtain k **where**

$k: Re\ A2 = k * Re\ A1\ Re\ B2 = k * Re\ B1\ Im\ B2 = k * Im\ B1$
using *linear-system-homogenous-3-2*[*of* $\lambda x\ y\ z. 1 * x + Re\ (u1 / u2) * y + Im\ (u1 / u2) * z\ 1\ Re\ (u1/u2)\ Im\ (u1/u2)$]

$$\lambda x\ y\ z. 1 * x + Re\ (v1 / v2) * y + Im\ (v1 / v2) * z\ 1\ Re\ (v1/v2)\ Im\ (v1/v2)$$

by (*auto simp add: field-simps*)

have $Re\ A2 \neq 0 \vee Re\ B2 \neq 0 \vee Im\ B2 \neq 0$
using $\langle ?H2 \neq mat-zero \rangle **$
by (*metis complex-cnj-zero complex-of-real-Re mat-zero-def of-real-0*)
hence $k \neq 0$
using k
by *auto*

show *circline-eq-cmat* $?H1\ ?H2$
using $**\ k\ \langle k \neq 0 \rangle$
by (*auto simp add: vec-cnj-def*) (*rule-tac* $x=k$ **in** *exI*, *auto simp add: complex.expand*)

qed

qed

qed

The only h-line that goes through zero and a non-zero point on the x-axis is the x-axis.

lemma *is-poincare-line-0-real-is-x-axis*:

assumes *is-poincare-line* $l\ 0_h \in \text{circline-set } l$
 $x \in \text{circline-set } l \cap \text{circline-set } x\text{-axis}\ x \neq 0_h\ x \neq \infty_h$
shows $l = x\text{-axis}$
using *assms*
using *is-poincare-line-trough-zero-trough-infty*[*OF* *assms*(1-2)]
using *unique-circline-set*[*of* $x\ 0_h\ \infty_h$]
by *auto*

The only h-line that goes through zero and a non-zero point on the y-axis is the y-axis.

lemma *is-poincare-line-0-imag-is-y-axis*:

assumes *is-poincare-line* $l\ 0_h \in \text{circline-set } l$
 $y \in \text{circline-set } l \cap \text{circline-set } y\text{-axis}\ y \neq 0_h\ y \neq \infty_h$
shows $l = y\text{-axis}$
using *assms*
using *is-poincare-line-trough-zero-trough-infty*[*OF* *assms*(1-2)]
using *unique-circline-set*[*of* $y\ 0_h\ \infty_h$]
by *auto*

4.1.6 H-isometries preserve h-lines

H-isometries are defined as homographies (actions of Möbius transformations) and antihomographies (compositions of actions of Möbius transformations with conjugation) that fix the unit disc (map it onto itself). They also map h-lines onto h-lines

We prove a bit more general lemma that states that all Möbius transformations that fix the unit circle (not necessarily the unit disc) map h-lines onto h-lines

```

lemma unit-circle-fix-preserve-is-poincare-line [simp]:
  assumes unit-circle-fix M is-poincare-line H
  shows is-poincare-line (moebius-circline M H)
  using assms
  unfolding is-poincare-line-iff
proof (safe)
  let ?H' = moebius-ocircline M (of-circline H)
  let ?U' = moebius-ocircline M ounit-circle
  assume ++: unit-circle-fix M perpendicular H unit-circle
  have ounit: ounit-circle = moebius-ocircline M ounit-circle  $\vee$ 
    ounit-circle = moebius-ocircline M (opposite-ocircline ounit-circle)
  using ++(1) unit-circle-fix-iff[of M]
  by (simp add: inj-of-ocircline moebius-circline-ocircline)

show perpendicular (moebius-circline M H) unit-circle
proof (cases pos-oriented ?H')
  case True
  hence *: of-circline (of-ocircline ?H') = ?H'
    using of-circline-of-ocircline-pos-oriented
    by blast
  from ounit show ?thesis
  proof
    assume **: ounit-circle = moebius-ocircline M ounit-circle
    show ?thesis
      using ++
      unfolding perpendicular-def
      by (simp, subst moebius-circline-ocircline, subst *, subst **) simp
  next
    assume **: ounit-circle = moebius-ocircline M (opposite-ocircline ounit-circle)
    show ?thesis
      using ++
      unfolding perpendicular-def
      by (simp, subst moebius-circline-ocircline, subst *, subst **) simp
  qed
next
  case False
  hence *: of-circline (of-ocircline ?H') = opposite-ocircline ?H'
    by (metis of-circline-of-ocircline pos-oriented-of-circline)
  from ounit show ?thesis
  proof
    assume **: ounit-circle = moebius-ocircline M ounit-circle
    show ?thesis
      using ++
      unfolding perpendicular-def
      by (simp, subst moebius-circline-ocircline, subst *, subst **) simp
  next
    assume **: ounit-circle = moebius-ocircline M (opposite-ocircline ounit-circle)
    show ?thesis
      using ++
      unfolding perpendicular-def
      by (simp, subst moebius-circline-ocircline, subst *, subst **) simp
  qed
qed
qed simp

```

```

lemma unit-circle-fix-preserve-is-poincare-line-iff [simp]:
  assumes unit-circle-fix M
  shows is-poincare-line (moebius-circline M H)  $\longleftrightarrow$  is-poincare-line H
  using assms
  using unit-circle-fix-preserve-is-poincare-line[of M H]
  using unit-circle-fix-preserve-is-poincare-line[of moebius-inv M moebius-circline M H]
  by (auto simp del: unit-circle-fix-preserve-is-poincare-line)

```

Since h-lines are preserved by transformations that fix the unit circle, so is collinearity.

```

lemma unit-disc-fix-preserve-poincare-collinear [simp]:
  assumes unit-circle-fix M poincare-collinear A

```

shows *poincare-collinear* (*moebius-pt* $M \text{ ' } A$)
 using *assms*
 unfolding *poincare-collinear-def*
 by (*auto*, *rule-tac* $x=\text{moebius-circline } M \text{ p in } exI$, *auto*)

lemma *unit-disc-fix-preserve-poincare-collinear-iff* [*simp*]:
 assumes *unit-circle-fix* M
 shows *poincare-collinear* (*moebius-pt* $M \text{ ' } A$) \longleftrightarrow *poincare-collinear* A
 using *assms*
 using *unit-disc-fix-preserve-poincare-collinear*[*of* $M \ A$]
 using *unit-disc-fix-preserve-poincare-collinear*[*of* *moebius-inv* M *moebius-pt* $M \text{ ' } A$]
 by (*auto simp del: unit-disc-fix-preserve-poincare-collinear*)

lemma *unit-disc-fix-preserve-poincare-collinear3* [*simp*]:
 assumes *unit-disc-fix* M
 shows *poincare-collinear* {*moebius-pt* $M \ u$, *moebius-pt* $M \ v$, *moebius-pt* $M \ w$ } \longleftrightarrow
 poincare-collinear { u , v , w }
 using *assms unit-disc-fix-preserve-poincare-collinear-iff*[*of* $M \ \{u, v, w\}$]
 by *simp*

Conjugation is also an h-isometry and it preserves h-lines.

lemma *is-poincare-line-conjugate-circline* [*simp*]:
 assumes *is-poincare-line* H
 shows *is-poincare-line* (*conjugate-circline* H)
 using *assms*
 by (*transfer*, *transfer*, *auto simp add: mat-cnj-def hermitean-def mat-adj-def*)

lemma *is-poincare-line-conjugate-circline-iff* [*simp*]:
 shows *is-poincare-line* (*conjugate-circline* H) \longleftrightarrow *is-poincare-line* H
 using *is-poincare-line-conjugate-circline*[*of* *conjugate-circline* H]
 by *auto*

Since h-lines are preserved by conjugation, so is collinearity.

lemma *conjugate-preserve-poincare-collinear* [*simp*]:
 assumes *poincare-collinear* A
 shows *poincare-collinear* (*conjugate* ' A)
 using *assms*
 unfolding *poincare-collinear-def*
 by *auto* (*rule-tac* $x=\text{conjugate-circline } p \text{ in } exI$, *auto*)

lemma *conjugate-conjugate* [*simp*]: *conjugate* ' *conjugate* ' $A = A$
 by (*auto simp add: image-iff*)

lemma *conjugate-preserve-poincare-collinear-iff* [*simp*]:
 shows *poincare-collinear* (*conjugate* ' A) \longleftrightarrow *poincare-collinear* A
 using *conjugate-preserve-poincare-collinear*[*of* A]
 using *conjugate-preserve-poincare-collinear*[*of* *conjugate* ' A]
 by (*auto simp del: conjugate-preserve-poincare-collinear*)

4.1.7 Mapping h-lines to x-axis

Each h-line in the Poincaré model can be mapped onto the x-axis (by a unit-disc preserving Möbius transformation).

lemma *ex-unit-disc-fix-is-poincare-line-to-x-axis*:
 assumes *is-poincare-line* l
 shows $\exists M. \text{unit-disc-fix } M \wedge \text{moebius-circline } M \ l = \text{x-axis}$
proof—
 from *assms* obtain $u \ v$ where $u \neq v \ u \in \text{unit-disc} \ v \in \text{unit-disc}$ and $\{u, v\} \subseteq \text{circline-set } l$
 using *ex-is-poincare-line-points*
 by *blast*
 then obtain M where $*$: *unit-disc-fix* M *moebius-pt* $M \ u = 0_h$ *moebius-pt* $M \ v \in \text{positive-x-axis}$
 using *ex-unit-disc-fix-to-zero-positive-x-axis*[*of* $u \ v$]
 by *auto*
moreover
 hence $\{0_h, \text{moebius-pt } M \ v\} \subseteq \text{circline-set } \text{x-axis}$


```

  unfolding positive-x-axis-def
  by auto
  moreover
  have moebius-pt M v ≠ 0h
    using ⟨u ≠ v⟩ *
    by (metis moebius-pt-neq-I)
  moreover
  have moebius-pt M v ≠ ∞h
    using ⟨unit-disc-fix M⟩ ⟨v ∈ unit-disc⟩
    using unit-disc-fix-discI
    by fastforce
  ultimately
  show ?thesis
    using ⟨is-poincare-line l⟩ ⟨{u, v} ⊆ circline-set l⟩ ⟨unit-disc-fix M⟩
    using is-poincare-line-0-real-is-x-axis[of moebius-circline M l moebius-pt M v]
    by (rule-tac x=M in exI, force)
qed

```

When proving facts about h-lines, without loss of generality it can be assumed that h-line is the x-axis (if the property being proved is invariant under Möbius transformations that fix the unit disc).

```

lemma wlog-line-x-axis:
  assumes is-line: is-poincare-line H
  assumes x-axis: P x-axis
  assumes preserves: ∧ M. [[unit-disc-fix M; P (moebius-circline M H)] ⇒ P H]
  shows P H
  using assms
  using ex-unit-disc-fix-is-poincare-line-to-x-axis[of H]
  by auto

```

4.2 Construction of the h-line between the two given points

Next we show how to construct the (unique) h-line between the two given points in the Poincaré model

Geometrically, h-line can be constructed by finding the inverse point of one of the two points and by constructing the circle (or line) through it and the two given points.

Algebraically, for two given points u and v in \mathbb{C} , the h-line matrix coefficients can be $A = i \cdot (u\bar{v} - v\bar{u})$ and $B = i \cdot (v(|u|^2 + 1) - u(|v|^2 + 1))$.

We need to extend this to homogenous coordinates. There are several degenerate cases.

- If $\{z, w\} = \{0_h, \infty_h\}$ then there is no unique h-line (any line through zero is an h-line).
- If z and w are mutually inverse, then the construction fails (both geometric and algebraic).
- If z and w are different points on the unit circle, then the standard construction fails (only geometric).
- None of this problematic cases occur when z and w are inside the unit disc.

We express the construction algebraically, and construct the Hermitean circline matrix for the two points given in homogenous coordinates. It works correctly in all cases except when the two points are the same or are mutually inverse.

```

definition mk-poincare-line-cmat :: real ⇒ complex ⇒ complex-mat where
  [simp]: mk-poincare-line-cmat A B = (cor A, B, cnj B, cor A)

```

```

lemma mk-poincare-line-cmat-zero-iff:
  mk-poincare-line-cmat A B = mat-zero ⟷ A = 0 ∧ B = 0
  by auto

```

```

lemma mk-poincare-line-cmat-hermitean
  [simp]: hermitean (mk-poincare-line-cmat A B)
  by simp

```

```

lemma mk-poincare-line-cmat-scale:
  cor k *sm mk-poincare-line-cmat A B = mk-poincare-line-cmat (k * A) (k * B)
  by simp

```

```

definition poincare-line-cvec-cmat :: complex-vec ⇒ complex-vec ⇒ complex-mat where
  [simp]: poincare-line-cvec-cmat z w =
    (let (z1, z2) = z;

```

```

(w1, w2) = w;
nom = w1*cnj w2*(z1*cnj z1 + z2*cnj z2) - z1*cnj z2*(w1*cnj w1 + w2*cnj w2);
den = z1*cnj z2*cnj w1*w2 - w1*cnj w2*cnj z1*z2
in if den ≠ 0 then
  mk-poincare-line-cmat (Re(i*den)) (i*nom)
else if z1*cnj z2 ≠ 0 then
  mk-poincare-line-cmat 0 (i*z1*cnj z2)
else if w1*cnj w2 ≠ 0 then
  mk-poincare-line-cmat 0 (i*w1*cnj w2)
else
  mk-poincare-line-cmat 0 i)

```

lemma *poincare-line-cvec-cmat-AeqD*:
assumes *poincare-line-cvec-cmat z w = (A, B, C, D)*
shows $A = D$
using *assms*
by (*cases z, cases w*) (*auto split: if-split-asm*)

lemma *poincare-line-cvec-cmat-hermitean [simp]*:
shows *hermitean (poincare-line-cvec-cmat z w)*
by (*cases z, cases w*) (*auto split: if-split-asm simp del: mk-poincare-line-cmat-def*)

lemma *poincare-line-cvec-cmat-nonzero [simp]*:
assumes $z \neq \text{vec-zero } w \neq \text{vec-zero}$
shows *poincare-line-cvec-cmat z w ≠ mat-zero*
proof –

obtain $z1\ z2\ w1\ w2$ **where** $*$: $z = (z1, z2)\ w = (w1, w2)$
by (*cases z, cases w, auto*)

let $?den = z1*cnj\ z2*cnj\ w1*w2 - w1*cnj\ w2*cnj\ z1*z2$

show *?thesis*

proof (*cases ?den ≠ 0*)

case *True*

have *is-real (i * ?den)*

using *eq-cn-j-iff-real[of i * ?den]*

by (*simp add: field-simps*)

hence $Re (i * ?den) \neq 0$

using $\langle ?den \neq 0 \rangle$

by (*metis complex-i-not-zero complex-surj mult-eq-0-iff zero-complex.code*)

thus *?thesis*

using $\langle ?den \neq 0 \rangle$

by (*simp del: mk-poincare-line-cmat-def mat-zero-def add: mk-poincare-line-cmat-zero-iff*)

next

case *False*

thus *?thesis*

using $*$

by (*simp del: mk-poincare-line-cmat-def mat-zero-def add: mk-poincare-line-cmat-zero-iff*)

qed

qed

lift-definition *poincare-line-hcoords-clmat :: complex-homo-coords ⇒ complex-homo-coords ⇒ circline-mat is poincare-line-cvec-cmat*
using *poincare-line-cvec-cmat-hermitean poincare-line-cvec-cmat-nonzero*
by *simp*

lift-definition *poincare-line :: complex-homo ⇒ complex-homo ⇒ circline is poincare-line-hcoords-clmat*

proof *transfer*

fix $z_a\ z_b\ w_a\ w_b$

assume $z_a \neq \text{vec-zero } z_b \neq \text{vec-zero } w_a \neq \text{vec-zero } w_b \neq \text{vec-zero}$

assume $z_a \approx_v z_b\ w_a \approx_v w_b$

obtain $z_{a1}\ z_{a2}\ z_{b1}\ z_{b2}\ w_{a1}\ w_{a2}\ w_{b1}\ w_{b2}$ **where**

$*$: $(z_{a1}, z_{a2}) = z_a\ (z_{b1}, z_{b2}) = z_b$

$(w_{a1}, w_{a2}) = w_a\ (w_{b1}, w_{b2}) = w_b$

by (*cases z_a, cases z_b, cases w_a, cases w_b, auto*)

obtain $k_z\ k_w$ **where**

$**$: $k_z \neq 0\ k_w \neq 0\ z_{b1} = k_z * z_{a1}\ z_{b2} = k_z * z_{a2}\ w_{b1} = k_w * w_{a1}\ w_{b2} = k_w * w_{a2}$

```

using ⟨za ≈v zb⟩ ⟨wa ≈v wb⟩ *[symmetric]
by auto

let ?nom = λ z1 z2 w1 w2. w1*cnj w2*(z1*cnj z1 + z2*cnj z2) - z1*cnj z2*(w1*cnj w1 + w2*cnj w2)
let ?den = λ z1 z2 w1 w2. z1*cnj z2*cnj w1*w2 - w1*cnj w2*cnj z1*z2

show circline-eq-cmat (poincare-line-cvec-cmat za wa)
  (poincare-line-cvec-cmat zb wb)
proof-
  have ∃ k. k ≠ 0 ∧
    poincare-line-cvec-cmat (zb1, zb2) (wb1, wb2) = cor k *sm poincare-line-cvec-cmat (za1, za2) (wa1, wa2)
proof (cases ?den za1 za2 wa1 wa2 ≠ 0)
  case True
  hence ?den zb1 zb2 wb1 wb2 ≠ 0
  using **
  by (simp add: field-simps)

  let ?k = kz * cnj kz * kw * cnj kw

  have ?k ≠ 0
  using **
  by simp

  have is-real ?k
  using eq-cnj-iff-real[of ?k]
  by auto

  have cor (Re ?k) = ?k
  using ⟨is-real ?k⟩
  using complex-of-real-Re
  by blast

  have Re ?k ≠ 0
  using ⟨?k ≠ 0⟩ ⟨cor (Re ?k) = ?k⟩
  by (metis of-real-0)

  have arg1: Re (i * ?den zb1 zb2 wb1 wb2) = Re ?k * Re (i * ?den za1 za2 wa1 wa2)
  apply (subst **)+
  apply (subst Re-mult-real[symmetric, OF ⟨is-real ?k⟩])
  apply (rule arg-cong[where f=Re])
  apply (simp add: field-simps)
  done
  have arg2: i * ?nom zb1 zb2 wb1 wb2 = ?k * i * ?nom za1 za2 wa1 wa2
  using **
  by (simp add: field-simps)
  have mk-poincare-line-cmat (Re (i*?den zb1 zb2 wb1 wb2)) (i*?nom zb1 zb2 wb1 wb2) =
    cor (Re ?k) *sm mk-poincare-line-cmat (Re (i*?den za1 za2 wa1 wa2)) (i*?nom za1 za2 wa1 wa2)
  using ⟨cor (Re ?k) = ?k⟩ ⟨is-real ?k⟩
  apply (subst mk-poincare-line-cmat-scale)
  apply (subst arg1, subst arg2)
  apply (subst ⟨cor (Re ?k) = ?k⟩)+
  apply simp
  done
  thus ?thesis
  using ⟨?den za1 za2 wa1 wa2 ≠ 0⟩ ⟨?den zb1 zb2 wb1 wb2 ≠ 0⟩
  using ⟨Re ?k ≠ 0⟩ ⟨cor (Re ?k) = ?k⟩
  by (rule-tac x=Re ?k in exI, simp)
next
  case False
  hence ?den zb1 zb2 wb1 wb2 = 0
  using **
  by (simp add: field-simps)
  show ?thesis
  proof (cases za1*cnj za2 ≠ 0)
  case True
  hence zb1*cnj zb2 ≠ 0

```

```

using **
by (simp add: field-simps)

let ?k = kz * cnj kz

have ?k ≠ 0 is-real ?k
  using **
  using eq-cnj-iff-real[of ?k]
  by auto
thus ?thesis
  using ⟨za1 * cnj za2 ≠ 0⟩ ⟨zb1 * cnj zb2 ≠ 0⟩
  using ⟨¬ (?den za1 za2 wa1 wa2 ≠ 0)⟩ ⟨?den zb1 zb2 wb1 wb2 = 0⟩ **
  by (rule-tac x=Re (kz * cnj kz) in exI, auto simp add: complex.expand)
next
case False
hence zb1 * cnj zb2 = 0
  using **
  by (simp add: field-simps)
show ?thesis
proof (cases wa1 * cnj wa2 ≠ 0)
case True
hence wb1 * cnj wb2 ≠ 0
  using **
  by (simp add: field-simps)

let ?k = kw * cnj kw

have ?k ≠ 0 is-real ?k
  using **
  using eq-cnj-iff-real[of ?k]
  by auto

thus ?thesis
  using ⟨¬ (za1 * cnj za2 ≠ 0)⟩
  using ⟨wa1 * cnj wa2 ≠ 0⟩ ⟨wb1 * cnj wb2 ≠ 0⟩
  using ⟨¬ (?den za1 za2 wa1 wa2 ≠ 0)⟩ ⟨?den zb1 zb2 wb1 wb2 = 0⟩ **
  by (rule-tac x=Re (kw * cnj kw) in exI)
  (auto simp add: complex.expand)
next
case False
hence wb1 * cnj wb2 = 0
  using **
  by (simp add: field-simps)
thus ?thesis
  using ⟨¬ (za1 * cnj za2 ≠ 0)⟩ ⟨zb1 * cnj zb2 = 0⟩
  using ⟨¬ (wa1 * cnj wa2 ≠ 0)⟩ ⟨wb1 * cnj wb2 = 0⟩
  using ⟨¬ (?den za1 za2 wa1 wa2 ≠ 0)⟩ ⟨?den zb1 zb2 wb1 wb2 = 0⟩ **
  by simp
qed
qed
qed
thus ?thesis
  using *[symmetric]
  by simp
qed
qed

```

4.2.1 Correctness of the construction

For finite points, our definition matches the classic algebraic definition for points in \mathbb{C} (given in ordinary, not homogenous coordinates).

lemma *poincare-line-non-homogenous:*

assumes $u \neq \infty_h$ $v \neq \infty_h$ $u \neq v$ $u \neq \text{inversion } v$

shows let $u' = \text{to-complex } u$; $v' = \text{to-complex } v$;

$$A = i * (u' * cnj v' - v' * cnj u');$$

$$B = i * (v' * ((cmod u')^2 + 1) - u' * ((cmod v')^2 + 1))$$

```

    in poincare-line u v = mk-circline A B (cnj B) A
using assms
unfolding unit-disc-def disc-def inversion-def
apply (simp add: Let-def)
proof (transfer, transfer, safe)
fix u1 u2 v1 v2
assume uv: (u1, u2) ≠ vec-zero (v1, v2) ≠ vec-zero
    ¬ (u1, u2) ≈v ∞v ¬ (v1, v2) ≈v ∞v
    ¬ (u1, u2) ≈v (v1, v2) ¬ (u1, u2) ≈v conjugate-cvec (reciprocal-cvec (v1, v2))
let ?u = to-complex-cvec (u1, u2) and ?v = to-complex-cvec (v1, v2)
let ?A = i * (?u * cnj ?v - ?v * cnj ?u)
let ?B = i * (?v * ((cor (cmod ?u))2 + 1) - ?u * ((cor (cmod ?v))2 + 1))
let ?C = -(i * (cnj ?v * ((cor (cmod ?u))2 + 1) - cnj ?u * ((cor (cmod ?v))2 + 1)))
let ?D = ?A
let ?H = (?A, ?B, ?C, ?D)

let ?den = u1 * cnj u2 * cnj v1 * v2 - v1 * cnj v2 * cnj u1 * u2

have u2 ≠ 0 v2 ≠ 0
    using uv
    using inf-cvec-z2-zero-iff
    by blast+

have ¬ (u1, u2) ≈v (cnj v2, cnj v1)
    using w(6)
    by (simp add: vec-cnj-def)
moreover
have (cnj v2, cnj v1) ≠ vec-zero
    using w(2)
    by auto
ultimately
have *: u1 * cnj v1 ≠ u2 * cnj v2 u1 * v2 ≠ u2 * v1
    using w(5) w(1) w(2) ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩
    using complex-cvec-eq-mix
    by blast+

show circline-eq-cmat (poincare-line-cvec-cmat (u1, u2) (v1, v2))
    (mk-circline-cmat ?A ?B ?C ?D)
proof (cases ?den ≠ 0)
    case True

    let ?nom = v1 * cnj v2 * (u1 * cnj u1 + u2 * cnj u2) - u1 * cnj u2 * (v1 * cnj v1 + v2 * cnj v2)
    let ?H' = mk-poincare-line-cmat (Re (i * ?den)) (i * ?nom)

    have circline-eq-cmat ?H ?H'
    proof–
    let ?k = (u2 * cnj v2) * (v2 * cnj u2)
    have is-real ?k
    using eq-cnj-iff-real
    by fastforce
    hence cor (Re ?k) = ?k
    using complex-of-real-Re
    by blast

    have Re (i * ?den) = Re ?k * ?A
    proof–
    have ?A = cnj ?A
    by (simp add: field-simps)
    hence is-real ?A
    using eq-cnj-iff-real
    by fastforce
    moreover
    have i * ?den = cnj (i * ?den)
    by (simp add: field-simps)
    hence is-real (i * ?den)

```

```

    using eq-cnj-iff-real
    by fastforce
  hence cor (Re (i * ?den)) = i * ?den
    using complex-of-real-Re
    by blast
  ultimately
  show ?thesis
    using ⟨cor (Re ?k) = ?k⟩
    by (simp add: field-simps)
qed

moreover
have i * ?nom = Re ?k * ?B
  using ⟨cor (Re ?k) = ?k⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ complex-mult-cnj-cmod[symmetric]
  by (auto simp add: field-simps)

moreover
have ?k ≠ 0
  using ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩
  by simp
hence Re ?k ≠ 0
  using ⟨is-real ?k⟩
  by (metis ⟨cor (Re ?k) = ?k⟩ of-real-0)

ultimately
show ?thesis
  by simp (rule-tac x=Re ?k in exI, simp add: mult.commute)
qed

moreover

have poincare-line-cvec-cmat (u1, u2) (v1, v2) = ?H'
  using ⟨?den ≠ 0⟩
  unfolding poincare-line-cvec-cmat-def
  by (simp add: Let-def)

moreover

hence hermitean ?H' ∧ ?H' ≠ mat-zero
  by (metis mk-poincare-line-cmat-hermitean poincare-line-cvec-cmat-nonzero uv(1) uv(2))

hence hermitean ?H ∧ ?H ≠ mat-zero
  using ⟨circline-eq-cmat ?H ?H'⟩
  using circline-eq-cmat-hermitean-nonzero[of ?H' ?H] symp-circline-eq-cmat
  unfolding symp-def
  by metis

hence mk-circline-cmat ?A ?B ?C ?D = ?H
  by simp

ultimately

have circline-eq-cmat (mk-circline-cmat ?A ?B ?C ?D)
  (poincare-line-cvec-cmat (u1, u2) (v1, v2))
  by simp
thus ?thesis
  using symp-circline-eq-cmat
  unfolding symp-def
  by blast
next
case False

let ?d = v1 * (u1 * cnj u1 / (u2 * cnj u2) + 1) / v2 - u1 * (v1 * cnj v1 / (v2 * cnj v2) + 1) / u2
let ?cd = cnj v1 * (u1 * cnj u1 / (u2 * cnj u2) + 1) / cnj v2 - cnj u1 * (v1 * cnj v1 / (v2 * cnj v2) + 1) / cnj
u2

```

```

have cnj ?d = ?cd
  by (simp add: mult.commute)

let ?d1 = (v1 / v2) * (cnj u1 / cnj u2) - 1
let ?d2 = u1 / u2 - v1 / v2

have **: ?d = ?d1 * ?d2
  using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩
  by (simp add: field-simps)

hence ?d ≠ 0
  using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ *
  by auto (simp add: field-simps)+

have is-real ?d1
proof-
  have cnj ?d1 = ?d1
    using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ *
    by (simp add: field-simps)
  thus ?thesis
    using eq-cnj-iff-real
    by blast
qed

show ?thesis
proof (cases u1 * cnj u2 ≠ 0)
case True
  let ?nom = u1 * cnj u2
  let ?H' = mk-poincare-line-cmat 0 (i * ?nom)

  have circline-eq-cmat ?H ?H'
  proof-

    let ?k = (u1 * cnj u2) / ?d

    have is-real ?k
    proof-
      have is-real ((u1 * cnj u2) / ?d2)
      proof-
        let ?rhs = (u2 * cnj u2) / (1 - (v1*u2)/(u1*v2))

        have 1: (u1 * cnj u2) / ?d2 = ?rhs
          using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ * ⟨u1 * cnj u2 ≠ 0⟩
          by (simp add: field-simps)
        moreover
        have cnj ?rhs = ?rhs
        proof-
          have cnj (1 - v1 * u2 / (u1 * v2)) = 1 - v1 * u2 / (u1 * v2)
            using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ * ⟨u1 * cnj u2 ≠ 0⟩
            by (simp add: field-simps)
          moreover
          have cnj (u2 * cnj u2) = u2 * cnj u2
            by simp
          ultimately
          show ?thesis
            by simp
        qed
      qed

    ultimately

    show ?thesis
      using eq-cnj-iff-real
      by fastforce
    qed

  thus ?thesis

```

```

    using ** ⟨is-real ?d1⟩
    by (metis complex-cnj-divide divide-divide-eq-left' eq-cnj-iff-real)
qed

have ?k ≠ 0
  using ⟨?d ≠ 0⟩ ⟨u1 * cnj u2 ≠ 0⟩
  by simp

have cnj ?k = ?k
  using ⟨is-real ?k⟩
  using eq-cnj-iff-real by blast

have Re ?k ≠ 0
  using ⟨?k ≠ 0⟩ ⟨is-real ?k⟩
  by (metis complex.expand zero-complex.simps(1) zero-complex.simps(2))

have u1 * cnj u2 = ?k * ?d
  using ⟨?d ≠ 0⟩
  by simp

moreover

hence cnj u1 * u2 = cnj ?k * cnj ?d
  by (metis complex-cnj-cnj complex-cnj-mult)
hence cnj u1 * u2 = ?k * ?cd
  using ⟨cnj ?k = ?k⟩ ⟨cnj ?d = ?cd⟩
  by metis

ultimately

show ?thesis
  using ⟨~ ?den ≠ 0⟩ ⟨u1 * cnj u2 ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ ⟨Re ?k ≠ 0⟩ ⟨is-real ?k⟩ ⟨?d ≠ 0⟩
  using complex-mult-cnj-cmod[symmetric, of u1]
  using complex-mult-cnj-cmod[symmetric, of v1]
  using complex-mult-cnj-cmod[symmetric, of u2]
  using complex-mult-cnj-cmod[symmetric, of v2]
  apply (simp add: power-divide norm-mult norm-divide)
  apply (rule-tac x=Re ?k in exI)
  apply simp
  apply (simp add: field-simps)
done
qed

moreover

have poincare-line-cvec-cmat (u1, u2) (v1, v2) = ?H'
  using ⟨¬ ?den ≠ 0⟩ ⟨u1 * cnj u2 ≠ 0⟩
  unfolding poincare-line-cvec-cmat-def
  by (simp add: Let-def)

moreover

hence hermitean ?H' ∧ ?H' ≠ mat-zero
  by (metis mk-poincare-line-cmat-hermitean poincare-line-cvec-cmat-nonzero uv(1) uv(2))

hence hermitean ?H ∧ ?H ≠ mat-zero
  using ⟨circline-eq-cmat ?H ?H'⟩
  using circline-eq-cmat-hermitean-nonzero[of ?H' ?H] symp-circline-eq-cmat
  unfolding symp-def
  by metis

hence mk-circline-cmat ?A ?B ?C ?D = ?H
  by simp

ultimately

```



```

have circline-eq-cmat (mk-circline-cmat ?A ?B ?C ?D)
  (poincare-line-cvec-cmat (u1, u2) (v1, v2))
  by simp
thus ?thesis
  using symp-circline-eq-cmat
  unfolding symp-def
  by blast
next
case False
show ?thesis
proof (cases v1 * cnj v2 ≠ 0)
  case True
  let ?nom = v1 * cnj v2
  let ?H' = mk-poincare-line-cmat 0 (i * ?nom)

  have circline-eq-cmat ?H ?H'
  proof -
    let ?k = (v1 * cnj v2) / ?d

    have is-real ?k
    proof -
      have is-real ((v1 * cnj v2) / ?d2)
      proof -
        let ?rhs = (v2 * cnj v2) / ((u1*v2)/(u2*v1) - 1)

        have 1: (v1 * cnj v2) / ?d2 = ?rhs
          using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ * ⟨v1 * cnj v2 ≠ 0⟩
          by (simp add: field-simps)
        moreover
        have cnj ?rhs = ?rhs
        proof -
          have cnj (u1 * v2 / (u2 * v1) - 1) = u1 * v2 / (u2 * v1) - 1
            using ⟨¬ ?den ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ * ⟨v1 * cnj v2 ≠ 0⟩
            by (simp add: field-simps)
          moreover
          have cnj (v2 * cnj v2) = v2 * cnj v2
            by simp
          ultimately
          show ?thesis
            by simp
        qed
      qed
    qed

    ultimately

  show ?thesis
    using eq-cnj-iff-real
    by fastforce
  qed

  thus ?thesis
    using ** ⟨is-real ?d1⟩
    by (metis complex-cnj-divide divide-divide-eq-left' eq-cnj-iff-real)
  qed

  have ?k ≠ 0
    using ⟨?d ≠ 0⟩ ⟨v1 * cnj v2 ≠ 0⟩
    by simp

  have cnj ?k = ?k
    using ⟨is-real ?k⟩
    using eq-cnj-iff-real by blast

  have Re ?k ≠ 0
    using ⟨?k ≠ 0⟩ ⟨is-real ?k⟩
    by (metis complex.expand zero-complex.simps(1) zero-complex.simps(2))

```

```

have v1 * cnj v2 = ?k * ?d
  using ⟨?d ≠ 0⟩
  by simp

moreover

hence cnj v1 * v2 = cnj ?k * cnj ?d
  by (metis complex-cnj-cnj complex-cnj-mult)
hence cnj v1 * v2 = ?k * ?cd
  using ⟨cnj ?k = ?k⟩ ⟨cnj ?d = ?cd⟩
  by metis

ultimately

show ?thesis
  using ⟨~ ?den ≠ 0⟩ ⟨v1 * cnj v2 ≠ 0⟩ ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩ ⟨Re ?k ≠ 0⟩ ⟨is-real ?k⟩ ⟨?d ≠ 0⟩
  using complex-mult-cnj-cmod[symmetric, of u1]
  using complex-mult-cnj-cmod[symmetric, of v1]
  using complex-mult-cnj-cmod[symmetric, of u2]
  using complex-mult-cnj-cmod[symmetric, of v2]
  apply (simp add: power-divide norm-mult norm-divide)
  apply (rule-tac x=Re ?k in ex1)
  apply simp
  apply (simp add: field-simps)
  done
qed

moreover

have poincare-line-cvec-cmat (u1, u2) (v1, v2) = ?H'
  using ⟨¬ ?den ≠ 0⟩ ⟨¬ u1 * cnj u2 ≠ 0⟩ ⟨v1 * cnj v2 ≠ 0⟩
  unfolding poincare-line-cvec-cmat-def
  by (simp add: Let-def)

moreover

hence hermitean ?H' ∧ ?H' ≠ mat-zero
  by (metis mk-poincare-line-cmat-hermitean poincare-line-cvec-cmat-nonzero uw(1) uw(2))

hence hermitean ?H ∧ ?H ≠ mat-zero
  using ⟨circline-eq-cmat ?H ?H'⟩
  using circline-eq-cmat-hermitean-nonzero[of ?H' ?H] symp-circline-eq-cmat
  unfolding symp-def
  by metis

hence mk-circline-cmat ?A ?B ?C ?D = ?H
  by simp

ultimately

have circline-eq-cmat (mk-circline-cmat ?A ?B ?C ?D)
  (poincare-line-cvec-cmat (u1, u2) (v1, v2))
  by simp
thus ?thesis
  using symp-circline-eq-cmat
  unfolding symp-def
  by blast
next
case False
hence False
  using ⟨¬ ?den ≠ 0⟩ ⟨¬ u1 * cnj u2 ≠ 0⟩ uw
  by (simp add: ⟨u2 ≠ 0⟩ ⟨v2 ≠ 0⟩)
thus ?thesis
  by simp
qed
qed

```

qed
qed

Our construction (in homogenous coordinates) always yields an h-line that contain two starting points (this also holds for all degenerate cases except when points are the same).

```

lemma poincare-line [simp]:
  assumes  $z \neq w$ 
  shows on-circline (poincare-line  $z$   $w$ )  $z$ 
         on-circline (poincare-line  $z$   $w$ )  $w$ 
proof-
  have on-circline (poincare-line  $z$   $w$ )  $z$   $\wedge$  on-circline (poincare-line  $z$   $w$ )  $w$ 
    using assms
  proof (transfer, transfer)
    fix  $z$   $w$ 
    assume vz:  $z \neq \text{vec-zero}$   $w \neq \text{vec-zero}$ 
    obtain  $z1$   $z2$   $w1$   $w2$  where
      zw:  $(z1, z2) = z$   $(w1, w2) = w$ 
      by (cases  $z$ , cases  $w$ , auto)

    let ?den =  $z1 * \text{cnj } z2 * \text{cnj } w1 * w2 - w1 * \text{cnj } w2 * \text{cnj } z1 * z2$ 
    have *:  $\text{cor } (\text{Re } (i * ?den)) = i * ?den$ 
    proof-
      have cnj ?den =  $-?den$ 
        by auto
      hence is-imag ?den
        using eq-minus-cnj-iff-imag[of ?den]
        by simp
      thus ?thesis
        using complex-of-real-Re[of i * ?den]
        by simp
    qed
  show on-circline-cmat-cvec (poincare-line-cvec-cmat  $z$   $w$ )  $z$   $\wedge$ 
        on-circline-cmat-cvec (poincare-line-cvec-cmat  $z$   $w$ )  $w$ 
    unfolding poincare-line-cvec-cmat-def mk-poincare-line-cmat-def
    apply (subst zw[symmetric])+
    unfolding Let-def prod.case
    apply (subst *)+
    by (auto simp add: vec-cnj-def field-simps)
  qed
  thus on-circline (poincare-line  $z$   $w$ )  $z$  on-circline (poincare-line  $z$   $w$ )  $w$ 
    by auto
qed

```

```

lemma poincare-line-circline-set [simp]:
  assumes  $z \neq w$ 
  shows  $z \in \text{circline-set } (\text{poincare-line } z \ w)$ 
          $w \in \text{circline-set } (\text{poincare-line } z \ w)$ 
  using assms
  by (auto simp add: circline-set-def)

```

When the points are different, the constructed line matrix always has a negative determinant

```

lemma poincare-line-type:
  assumes  $z \neq w$ 
  shows circline-type (poincare-line  $z$   $w$ ) =  $-1$ 
proof-
  have  $\exists a \ b. a \neq b \wedge \{a, b\} \subseteq \text{circline-set } (\text{poincare-line } z \ w)$ 
    using poincare-line[of z w] assms
    unfolding circline-set-def
    by (rule-tac x=z in exI, rule-tac x=w in exI, simp)
  thus ?thesis
    using circline-type[of poincare-line z w]
    using circline-type-pos-card-eq0[of poincare-line z w]
    using circline-type-zero-card-eq1[of poincare-line z w]
    by auto
qed

```

The constructed line is an h-line in the Poincaré model (in all cases when the two points are different)

```

lemma is-poincare-line-poincare-line [simp]:
  assumes  $z \neq w$ 
  shows is-poincare-line (poincare-line  $z w$ )
  using poincare-line-type[of  $z w$ , OF assms]
proof (transfer, transfer)
  fix  $z w$ 
  assume  $vz: z \neq \text{vec-zero } w \neq \text{vec-zero}$ 
  obtain  $A B C D$  where  $*$ : poincare-line-cvec-cmat  $z w = (A, B, C, D)$ 
    by (cases poincare-line-cvec-cmat  $z w$ ) auto
  assume circline-type-cmat (poincare-line-cvec-cmat  $z w$ ) = - 1
  thus is-poincare-line-cmat (poincare-line-cvec-cmat  $z w$ )
    using  $vz *$ 
    using poincare-line-cvec-cmat-hermitean[of  $z w$ ]
    using poincare-line-cvec-cmat-nonzero[of  $z w$ ]
    using poincare-line-cvec-cmat-AeqD[of  $z w A B C D$ ]
    using hermitean-elems[of  $A B C D$ ]
    using cmod-power2[of  $D$ ] cmod-power2[of  $C$ ]
    unfolding is-poincare-line-cmat-def
    by (simp del: poincare-line-cvec-cmat-def add: sgn-1-neg power2-eq-square)
qed

```

When the points are different, the constructed h-line between two points also contains their inverses

```

lemma poincare-line-inversion:
  assumes  $z \neq w$ 
  shows on-circline (poincare-line  $z w$ ) (inversion  $z$ )
    on-circline (poincare-line  $z w$ ) (inversion  $w$ )
  using assms
  using is-poincare-line-poincare-line[OF  $\langle z \neq w \rangle$ ]
  using is-poincare-line-inverse-point
  unfolding circline-set-def
  by auto

```

When the points are different, the onstructed h-line between two points contains the inverse of its every point

```

lemma poincare-line-inversion-full:
  assumes  $u \neq v$ 
  assumes on-circline (poincare-line  $u v$ )  $x$ 
  shows on-circline (poincare-line  $u v$ ) (inversion  $x$ )
  using is-poincare-line-inverse-point[of poincare-line  $u v x$ ]
  using is-poincare-line-poincare-line[OF  $\langle u \neq v \rangle$ ] assms
  unfolding circline-set-def
  by simp

```

4.2.2 Existence of h-lines

There is an h-line trough every point in the Poincaré model

```

lemma ex-poincare-line-one-point:
  shows  $\exists l. \text{is-poincare-line } l \wedge z \in \text{circline-set } l$ 
proof (cases  $z = 0_h$ )
  case True
  thus ?thesis
    by (rule-tac  $x=x\text{-axis}$  in exI) simp
next
  case False
  thus ?thesis
    by (rule-tac  $x=\text{poincare-line } 0_h z$  in exI) auto
qed

```

```

lemma poincare-collinear-singleton [simp]:
  assumes  $u \in \text{unit-disc}$ 
  shows poincare-collinear  $\{u\}$ 
  using assms
  using ex-poincare-line-one-point[of  $u$ ]
  by (auto simp add: poincare-collinear-def)

```

There is an h-line through every two points in the Poincaré model

lemma *ex-poincare-line-two-points*:
assumes $z \neq w$
shows $\exists l. \text{is-poincare-line } l \wedge z \in \text{circline-set } l \wedge w \in \text{circline-set } l$
using *assms*
by (*rule-tac x=poincare-line z w in exI, simp*)

lemma *poincare-collinear-doubleton* [*simp*]:
assumes $u \in \text{unit-disc } v \in \text{unit-disc}$
shows *poincare-collinear* $\{u, v\}$
using *assms*
using *ex-poincare-line-one-point*[*of u*]
using *ex-poincare-line-two-points*[*of u v*]
by (*cases u = v*) (*simp-all add: poincare-collinear-def*)

4.2.3 Uniqueness of h-lines

The only h-line between two points is the one obtained by the line-construction.

First we show this only for two different points inside the disc.

lemma *unique-poincare-line*:
assumes *in-disc*: $u \neq v \ u \in \text{unit-disc } v \in \text{unit-disc}$
assumes *on-l*: $u \in \text{circline-set } l \ v \in \text{circline-set } l \ \text{is-poincare-line } l$
shows $l = \text{poincare-line } u \ v$
using *assms*
using *unique-is-poincare-line*[*of u v l poincare-line u v*]
unfolding *circline-set-def*
by *auto*

The assumption that the points are inside the disc can be relaxed.

lemma *unique-poincare-line-general*:
assumes *in-disc*: $u \neq v \ u \neq \text{inversion } v$
assumes *on-l*: $u \in \text{circline-set } l \ v \in \text{circline-set } l \ \text{is-poincare-line } l$
shows $l = \text{poincare-line } u \ v$
using *assms*
using *unique-is-poincare-line-general*[*of u v l poincare-line u v*]
unfolding *circline-set-def*
by *auto*

The explicit line construction enables us to prove that there exists a unique h-line through any given two h-points (uniqueness part was already shown earlier).

First we show this only for two different points inside the disc.

lemma *ex1-poincare-line*:
assumes $u \neq v \ u \in \text{unit-disc } v \in \text{unit-disc}$
shows $\exists! l. \text{is-poincare-line } l \wedge u \in \text{circline-set } l \wedge v \in \text{circline-set } l$
proof (*rule ex1I*)
let $?l = \text{poincare-line } u \ v$
show $\text{is-poincare-line } ?l \wedge u \in \text{circline-set } ?l \wedge v \in \text{circline-set } ?l$
using *assms*
unfolding *circline-set-def*
by *auto*
next
fix l
assume $\text{is-poincare-line } l \wedge u \in \text{circline-set } l \wedge v \in \text{circline-set } l$
thus $l = \text{poincare-line } u \ v$
using *unique-poincare-line assms*
by *auto*
qed

The assumption that the points are in the disc can be relaxed.

lemma *ex1-poincare-line-general*:
assumes $u \neq v \ u \neq \text{inversion } v$
shows $\exists! l. \text{is-poincare-line } l \wedge u \in \text{circline-set } l \wedge v \in \text{circline-set } l$
proof (*rule ex1I*)

```

let ?l = poincare-line u v
show is-poincare-line ?l  $\wedge$  u  $\in$  circline-set ?l  $\wedge$  v  $\in$  circline-set ?l
  using assms
  unfolding circline-set-def
  by auto
next
fix l
assume is-poincare-line l  $\wedge$  u  $\in$  circline-set l  $\wedge$  v  $\in$  circline-set l
thus l = poincare-line u v
  using unique-poincare-line-general assms
  by auto
qed

```

4.2.4 Some consequences of line uniqueness

H-line uv is the same as the h-line vu .

```

lemma poincare-line-sym:
  assumes u  $\in$  unit-disc v  $\in$  unit-disc u  $\neq$  v
  shows poincare-line u v = poincare-line v u
  using assms
  using unique-poincare-line[of u v poincare-line v u]
  by simp

```

```

lemma poincare-line-sym-general:
  assumes u  $\neq$  v u  $\neq$  inversion v
  shows poincare-line u v = poincare-line v u
  using assms
  using unique-poincare-line-general[of u v poincare-line v u]
  by simp

```

Each h-line is the h-line constructed out of its two arbitrary different points.

```

lemma ex-poincare-line-points:
  assumes is-poincare-line H
  shows  $\exists$  u v. u  $\in$  unit-disc  $\wedge$  v  $\in$  unit-disc  $\wedge$  u  $\neq$  v  $\wedge$  H = poincare-line u v
  using assms
  using ex-is-poincare-line-points
  using unique-poincare-line[where l=H]
  by fastforce

```

If an h-line contains two different points on x-axis/y-axis then it is the x-axis/y-axis.

```

lemma poincare-line-0-real-is-x-axis:
  assumes x  $\in$  circline-set x-axis x  $\neq$  0h x  $\neq$   $\infty$ h
  shows poincare-line 0h x = x-axis
  using assms
  using is-poincare-line-0-real-is-x-axis[of poincare-line 0h x x]
  by auto

```

```

lemma poincare-line-0-imag-is-y-axis:
  assumes y  $\in$  circline-set y-axis y  $\neq$  0h y  $\neq$   $\infty$ h
  shows poincare-line 0h y = y-axis
  using assms
  using is-poincare-line-0-imag-is-y-axis[of poincare-line 0h y y]
  by auto

```

```

lemma poincare-line-x-axis:
  assumes x  $\in$  unit-disc y  $\in$  unit-disc x  $\in$  circline-set x-axis y  $\in$  circline-set x-axis x  $\neq$  y
  shows poincare-line x y = x-axis
  using assms
  using unique-poincare-line
  by auto

```

```

lemma poincare-line-minus-one-one [simp]:
  shows poincare-line (of-complex (-1)) (of-complex 1) = x-axis
proof -
  have 0h  $\in$  circline-set (poincare-line (of-complex (-1)) (of-complex 1))

```

unfolding *circline-set-def*
by *simp (transfer, transfer, simp add: vec-cnj-def)*
hence *poincare-line 0_h (of-complex 1) = poincare-line (of-complex (-1)) (of-complex 1)*
by (*metis is-poincare-line-poincare-line is-poincare-line-trough-zero-trough-infty not-zero-on-unit-circle of-complex-inj of-complex-one one-neq-neg-one one-on-unit-circle poincare-line-0-real-is-x-axis poincare-line-circline-set(2) reciprocal-involution reciprocal-one reciprocal-zero unique-circline-01inf*)
thus *?thesis*
using *poincare-line-0-real-is-x-axis[of of-complex 1]*
by *auto*
qed

4.2.5 Transformations of constructed lines

Unit discs preserving Möbius transformations preserve the h-line construction

lemma *unit-disc-fix-preserve-poincare-line [simp]:*
assumes *unit-disc-fix M u ∈ unit-disc v ∈ unit-disc u ≠ v*
shows *poincare-line (moebius-pt M u) (moebius-pt M v) = moebius-circline M (poincare-line u v)*
proof (*rule unique-poincare-line[symmetric]*)
show *moebius-pt M u ≠ moebius-pt M v*
using *⟨u ≠ v⟩*
by *auto*
next
show *moebius-pt M u ∈ circline-set (moebius-circline M (poincare-line u v))*
moebius-pt M v ∈ circline-set (moebius-circline M (poincare-line u v))
unfolding *circline-set-def*
using *moebius-circline[of M poincare-line u v] ⟨u ≠ v⟩*
by *auto*
next
from *assms(1)* **have** *unit-circle-fix M*
by *simp*
thus *is-poincare-line (moebius-circline M (poincare-line u v))*
using *unit-circle-fix-preserve-is-poincare-line assms*
by *auto*
next
show *moebius-pt M u ∈ unit-disc moebius-pt M v ∈ unit-disc*
using *assms(2-3) unit-disc-fix-iff[OF assms(1)]*
by *auto*
qed

Conjugate preserve the h-line construction

lemma *conjugate-preserve-poincare-line [simp]:*
assumes *u ∈ unit-disc v ∈ unit-disc u ≠ v*
shows *poincare-line (conjugate u) (conjugate v) = conjugate-circline (poincare-line u v)*
proof –
have *conjugate u ≠ conjugate v*
using *⟨u ≠ v⟩*
by (*auto simp add: conjugate-inj*)
moreover
have *conjugate u ∈ unit-disc conjugate v ∈ unit-disc*
using *assms*
by *auto*
moreover
have *conjugate u ∈ circline-set (conjugate-circline (poincare-line u v))*
conjugate v ∈ circline-set (conjugate-circline (poincare-line u v))
using *⟨u ≠ v⟩*
by *simp-all*
moreover
have *is-poincare-line (conjugate-circline (poincare-line u v))*
using *is-poincare-line-poincare-line[OF ⟨u ≠ v⟩]*
by *simp*
ultimately
show *?thesis*
using *unique-poincare-line[of conjugate u conjugate v conjugate-circline (poincare-line u v)]*
by *simp*
qed

4.2.6 Collinear points and h-lines

lemma *poincare-collinear3-poincare-line-general*:
assumes *poincare-collinear* {*a*, *a1*, *a2*} *a1* \neq *a2* *a1* \neq *inversion a2*
shows *a* \in *circline-set* (*poincare-line a1 a2*)
using *assms*
using *poincare-collinear-def unique-poincare-line-general*
by *auto*

lemma *poincare-line-poincare-collinear3-general*:
assumes *a* \in *circline-set* (*poincare-line a1 a2*) *a1* \neq *a2*
shows *poincare-collinear* {*a*, *a1*, *a2*}
using *assms*
unfolding *poincare-collinear-def*
by (*rule-tac x=poincare-line a1 a2 in exI, simp*)

lemma *poincare-collinear3-poincare-lines-equal-general*:
assumes *poincare-collinear* {*a*, *a1*, *a2*} *a* \neq *a1* *a* \neq *a2* *a* \neq *inversion a1* *a* \neq *inversion a2*
shows *poincare-line a a1* = *poincare-line a a2*
using *assms*
using *unique-poincare-line-general[of a a2 poincare-line a a1]*
by (*simp add: insert-commute poincare-collinear3-poincare-line-general*)

4.2.7 Points collinear with 0_h

lemma *poincare-collinear-zero-iff*:
assumes *of-complex y' \in unit-disc* **and** *of-complex z' \in unit-disc* **and**
 $y' \neq z'$ **and** $y' \neq 0$ **and** $z' \neq 0$
shows *poincare-collinear* { 0_h , *of-complex y'*, *of-complex z'*} \longleftrightarrow
 $y' * cnj z' = cnj y' * z'$ (**is** *?lhs \longleftrightarrow ?rhs*)

proof –

have *of-complex y' \neq of-complex z'*
using *assms*
using *of-complex-inj*
by *blast*

show *?thesis*

proof

assume *?lhs*

hence $0_h \in$ *circline-set* (*poincare-line* (*of-complex y'*) (*of-complex z'*))
using *unique-poincare-line[of of-complex y' of-complex z']*
using *assms* \langle *of-complex y' \neq of-complex z'* \rangle
unfolding *poincare-collinear-def*
by *auto*

moreover

let *?mix* = $y' * cnj z' - cnj y' * z'$

have *is-real* (*i * ?mix*)

using *eq-cnj-iff-real[of ?mix]*

by *auto*

hence $y' * cnj z' = cnj y' * z' \longleftrightarrow$ Re (*i * ?mix*) = 0

using *complex.expand[of i * ?mix 0]*

by (*metis complex-i-not-zero eq-iff-diff-eq-0 mult-eq-0-iff zero-complex.simps(1) zero-complex.simps(2)*)

ultimately

show *?rhs*

using $\langle y' \neq z' \rangle$ $\langle y' \neq 0 \rangle$ $\langle z' \neq 0 \rangle$

unfolding *circline-set-def*

by *simp* (*transfer, transfer, auto simp add: vec-cnj-def split: if-split-asm, metis Re-complex-of-real Re-mult-real*

Im-complex-of-real)

next

assume *?rhs*

thus *?lhs*

using *assms* \langle *of-complex y' \neq of-complex z'* \rangle

unfolding *poincare-collinear-def*

unfolding *circline-set-def*

apply (*rule-tac x=poincare-line* (*of-complex y'*) (*of-complex z'*) **in** *exI*)

apply *auto*

apply (*transfer, transfer, simp add: vec-cnj-def*)


```

done
qed
qed

lemma poincare-collinear-zero-polar-form:
  assumes poincare-collinear {0h, of-complex x, of-complex y} and
    x ≠ 0 and y ≠ 0 and of-complex x ∈ unit-disc and of-complex y ∈ unit-disc
  shows ∃ φ rx ry. x = cor rx * cis φ ∧ y = cor ry * cis φ ∧ rx ≠ 0 ∧ ry ≠ 0
proof-
  from ⟨x ≠ 0⟩ ⟨y ≠ 0⟩ obtain φ φ' rx ry where
    polar: x = cor rx * cis φ y = cor ry * cis φ' and φ = Arg x φ' = Arg y
  by (metis cmod-cis)
  hence rx ≠ 0 ry ≠ 0
  using ⟨x ≠ 0⟩ ⟨y ≠ 0⟩
  by auto
  have of-complex y ∈ circline-set (poincare-line 0h (of-complex x))
  using assms
  using unique-poincare-line[of 0h of-complex x]
  unfolding poincare-collinear-def
  unfolding circline-set-def
  using of-complex-zero-iff
  by fastforce
  hence cnj x * y = x * cnj y
  using ⟨x ≠ 0⟩ ⟨y ≠ 0⟩
  unfolding circline-set-def
  by simp (transfer, transfer, simp add: vec-cnj-def field-simps)
  hence cis(φ' - φ) = cis(φ - φ')
  using polar ⟨rx ≠ 0⟩ ⟨ry ≠ 0⟩
  by (simp add: cis-mult)
  hence sin (φ' - φ) = 0
  using cis-diff-cis-opposite[of φ' - φ]
  by simp
  then obtain k :: int where φ' - φ = k * pi
  using sin-zero-iff-int2[of φ' - φ]
  by auto
  hence *: φ' = φ + k * pi
  by simp
  show ?thesis
  proof (cases even k)
    case True
    then obtain k' where k = 2*k'
    using evenE by blast
    hence cis φ = cis φ'
    using * cos-periodic-int sin-periodic-int
    by (simp add: cis.ctr field-simps)
    thus ?thesis
    using polar ⟨rx ≠ 0⟩ ⟨ry ≠ 0⟩
    by (rule-tac x=φ in exI, rule-tac x=rx in exI, rule-tac x=ry in exI) simp
  next
    case False
    then obtain k' where k = 2*k' + 1
    using oddE by blast
    hence cis φ = - cis φ'
    using * cos-periodic-int sin-periodic-int
    by (simp add: cis.ctr complex-minus field-simps)
    thus ?thesis
    using polar ⟨rx ≠ 0⟩ ⟨ry ≠ 0⟩
    by (rule-tac x=φ in exI, rule-tac x=rx in exI, rule-tac x=-ry in exI) simp
  qed
qed
qed

end
theory Poincare-Lines-Ideal-Points
imports Poincare-Lines
begin

```

4.3 Ideal points of h-lines

Ideal points of an h-line are points where the h-line intersects the unit disc.

4.3.1 Calculation of ideal points

We decided to define ideal points constructively, i.e., we calculate the coordinates of ideal points for a given h-line explicitly. Namely, if the h-line is determined by A and B , the two intersection points are

$$\frac{B}{|B|^2} \left(-A \pm i \cdot \sqrt{|B|^2 - A^2} \right).$$

definition *calc-ideal-point1-cvec* :: *complex* \Rightarrow *complex* \Rightarrow *complex-vec* **where**

[simp]: *calc-ideal-point1-cvec* $A B =$
 (let *discr* = $\text{Re} ((\text{cmod } B)^2 - (\text{Re } A)^2)$ in
 ($B * (-A - i * \text{sqrt}(\text{discr}))$, $(\text{cmod } B)^2$))

definition *calc-ideal-point2-cvec* :: *complex* \Rightarrow *complex* \Rightarrow *complex-vec* **where**

[simp]: *calc-ideal-point2-cvec* $A B =$
 (let *discr* = $\text{Re} ((\text{cmod } B)^2 - (\text{Re } A)^2)$ in
 ($B * (-A + i * \text{sqrt}(\text{discr}))$, $(\text{cmod } B)^2$))

definition *calc-ideal-points-cmat-cvec* :: *complex-mat* \Rightarrow *complex-vec set* **where**

[simp]: *calc-ideal-points-cmat-cvec* $H =$
 (if *is-poincare-line-cmat* H then
 let $(A, B, C, D) = H$
 in {*calc-ideal-point1-cvec* $A B$, *calc-ideal-point2-cvec* $A B$ }
 else
 $\{(-1, 1), (1, 1)\}$)

lift-definition *calc-ideal-points-clmat-hcoords* :: *circline-mat* \Rightarrow *complex-homo-coords set* **is** *calc-ideal-points-cmat-cvec*
by (*auto simp add: Let-def split: if-split-asm*)

lift-definition *calc-ideal-points* :: *circline* \Rightarrow *complex-homo set* **is** *calc-ideal-points-clmat-hcoords*

proof *transfer*

fix $H1 H2$

assume *hh*: *hermitean* $H1 \wedge H1 \neq \text{mat-zero}$ *hermitean* $H2 \wedge H2 \neq \text{mat-zero}$

obtain $A1 B1 C1 D1 A2 B2 C2 D2$ **where** $*$: $H1 = (A1, B1, C1, D1)$ $H2 = (A2, B2, C2, D2)$

by (*cases* $H1$, *cases* $H2$, *auto*)

assume *circline-eq-cmat* $H1 H2$

then obtain k **where** $k: k \neq 0$ $H2 = \text{cor } k *_{sm} H1$

by *auto*

thus *rel-set* (\approx_v) (*calc-ideal-points-cmat-cvec* $H1$) (*calc-ideal-points-cmat-cvec* $H2$)

proof (*cases is-poincare-line-cmat* $H1$)

case *True*

hence *is-poincare-line-cmat* $H2$

using $k * \text{hermitean-mult-real}$ [of $H1$ k] *hh*

by (*auto simp add: power2-eq-square norm-mult*)

have $*$: $\text{sqrt} (|k| * \text{cmod } B1 * (|k| * \text{cmod } B1) - k * \text{Re } D1 * (k * \text{Re } D1)) =$
 $|k| * \text{sqrt}(\text{cmod } B1 * \text{cmod } B1 - \text{Re } D1 * \text{Re } D1)$

proof—

have $|k| * \text{cmod } B1 * (|k| * \text{cmod } B1) - k * \text{Re } D1 * (k * \text{Re } D1) =$
 $k^2 * (\text{cmod } B1 * \text{cmod } B1 - \text{Re } D1 * \text{Re } D1)$

by (*simp add: power2-eq-square field-simps*)

thus *?thesis*

by (*simp add: real-sqrt-mult*)

qed

show *?thesis*

using $\langle \text{is-poincare-line-cmat } H1 \rangle \langle \text{is-poincare-line-cmat } H2 \rangle$

using $*$ k

apply (*simp add: Let-def*)

apply *safe*

apply (*simp add: power2-eq-square rel-set-def norm-mult*)

apply *safe*

apply (*cases* $k > 0$)

apply (*rule-tac* $x=(\text{cor } k)^2$ **in** *exI*)

```

    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (erule notE, rule-tac x=(cor k)2 in exI)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (cases k > 0)
    apply (erule notE, rule-tac x=(cor k)2 in exI)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (rule-tac x=(cor k)2 in exI)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (cases k > 0)
    apply (rule-tac x=(cor k)2 in exI)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (erule notE, rule-tac x=(cor k)2 in exI)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (rule-tac x=(cor k)2 in exI)
    apply (cases k > 0)
    apply (erule notE, rule-tac x=(cor k)2 in exI)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    apply (subst **)
    apply (simp add: power2-eq-square field-simps)
    done
next
case False
hence ¬ is-poincare-line-cmat H2
using k * hermitean-mult-real[of H1 k] hh
by (auto simp add: power2-eq-square norm-mult)
have rel-set (≈v) {(- 1, 1), (1, 1)} {(- 1, 1), (1, 1)}
by (simp add: rel-set-def)
thus ?thesis
using ⟨¬ is-poincare-line-cmat H1⟩ ⟨¬ is-poincare-line-cmat H2⟩
using *
by (auto simp add: Let-def)
qed
qed

```

Correctness of the calculation

We show that for every h-line its two calculated ideal points are different and are on the intersection of that line and the unit circle.

Calculated ideal points are on the unit circle

lemma *calc-ideal-point-1-unit:*

assumes *is-real A (cmod B)² > (Re A)²*

assumes *(z1, z2) = calc-ideal-point1-cvec A B*

shows *z1 * cnj z1 = z2 * cnj z2*

proof—

let *?discr = Re ((cmod B)² - (Re A)²)*

have *?discr > 0*

using *assms*

by *(simp add: cmod-power2)*

have *(B*(-A - i*sqrt(?discr))) * cnj (B*(-A - i*sqrt(?discr))) = (B * cnj B) * (A² + cor (abs ?discr))*

using *⟨is-real A⟩ eq-cnj-iff-real[of A]*

by *(simp add: field-simps power2-eq-square)*

also have *... = (B * cnj B) * (cmod B)²*

using *⟨?discr > 0⟩*

using *assms*

using *complex-of-real-Re[of (cmod B)² - (Re A)²] complex-of-real-Re[of A] ⟨is-real A⟩*

by *(simp add: power2-eq-square)*

also have *... = (cmod B)² * cnj ((cmod B)²)*

using *complex-cnj-complex-of-real complex-mult-cnj-cmod*

by *presburger*
finally show *?thesis*
 using *assms*
 by *simp*
qed

lemma *calc-ideal-point-2-unit:*

assumes *is-real A (cmod B)² > (cmod A)²*
assumes *(z1, z2) = calc-ideal-point2-cvec A B*
shows *z1 * cnj z1 = z2 * cnj z2*

proof–

let *?discr = Re ((cmod B)² - (Re A)²)*
have *?discr > 0*
 using *assms*
 by (*simp add: cmod-power2*)
have *(B*(-A + i*sqrt(?discr))) * cnj (B*(-A + i*sqrt(?discr))) = (B * cnj B) * (A² + cor (abs ?discr))*
 using *<is-real A> eq-cnj-iff-real[of A]*
 by (*simp add: field-simps power2-eq-square*)
also have *... = (B * cnj B) * (cmod B)²*
 using *<?discr > 0>*
 using *assms*
 using *complex-of-real-Re[of (cmod B)² - (Re A)²] complex-of-real-Re[of A] <is-real A>*
 by (*simp add: power2-eq-square*)
also have *... = (cmod B)² * cnj ((cmod B)²)*
 using *complex-cnj-complex-of-real complex-mult-cnj-cmod*
 by *presburger*
finally show *?thesis*
 using *assms*
 by *simp*

qed

lemma *calc-ideal-points-on-unit-circle:*

shows $\forall z \in \text{calc-ideal-points } H. z \in \text{circline-set unit-circle}$
unfolding *circline-set-def*
apply *simp*

proof (*transfer, transfer*)

fix *H*
assume *hh: hermitean H ∧ H ≠ mat-zero*
obtain *A B C D* **where** **: H = (A, B, C, D)*
 by (*cases H, auto*)
have $\forall (z1, z2) \in \text{calc-ideal-points-cmat-cvec } H. z1 * cnj z1 = z2 * cnj z2$
 using *hermitean-elems[of A B C D]*
 unfolding *calc-ideal-points-cmat-cvec-def*
 using *calc-ideal-point-1-unit[of A B]*
 using *calc-ideal-point-2-unit[of A B]*
 using *hh **
apply (*cases calc-ideal-point1-cvec A B, cases calc-ideal-point2-cvec A B*)
apply (*auto simp add: Let-def simp del: calc-ideal-point1-cvec-def calc-ideal-point2-cvec-def*)
done
thus *Ball (calc-ideal-points-cmat-cvec H) (on-circline-cmat-cvec unit-circle-cmat)*
 using *on-circline-cmat-cvec-unit*
 by (*auto simp del: on-circline-cmat-cvec-def calc-ideal-points-cmat-cvec-def*)

qed

Calculated ideal points are on the h-line

lemma *calc-ideal-point1-sq:*

assumes *(z1, z2) = calc-ideal-point1-cvec A B is-real A (cmod B)² > (cmod A)²*
shows *z1 * cnj z1 + z2 * cnj z2 = 2 * (B * cnj B)²*

proof–

let *?discr = Re ((cmod B)² - (Re A)²)*
have *?discr > 0*
 using *assms*
 by (*simp add: cmod-power2*)
have *z1 * cnj z1 = (B * cnj B) * (-A + i*sqrt(?discr))*(-A - i*sqrt(?discr))*
 using *assms eq-cnj-iff-real[of A]*
 by (*simp*)

also have $\dots = (B * \text{cnj } B) * (A^2 + ?discr)$
using *complex-of-real-Re*[of A] $\langle is\text{-real } A \rangle \langle ?discr > 0 \rangle$
by (*simp add: power2-eq-square field-simps*)
finally
have $z1 * \text{cnj } z1 = (B * \text{cnj } B)^2$
using *complex-of-real-Re*[of $(\text{cmod } B)^2 - (\text{Re } A)^2$] *complex-of-real-Re*[of A] $\langle is\text{-real } A \rangle$
using *complex-mult-cnj-cmod*[of B]
by (*simp add: power2-eq-square*)
moreover
have $z2 * \text{cnj } z2 = (B * \text{cnj } B)^2$
using *assms*
by *simp*
ultimately
show *?thesis*
by *simp*
qed

lemma *calc-ideal-point2-sq*:

assumes $(z1, z2) = \text{calc-ideal-point2-cvec } A \ B \ is\text{-real } A \ (\text{cmod } B)^2 > (\text{cmod } A)^2$
shows $z1 * \text{cnj } z1 + z2 * \text{cnj } z2 = 2 * (B * \text{cnj } B)^2$

proof–

let $?discr = \text{Re } ((\text{cmod } B)^2 - (\text{Re } A)^2)$
have $?discr > 0$
using *assms*
by (*simp add: cmod-power2*)
have $z1 * \text{cnj } z1 = (B * \text{cnj } B) * (-A + i*\text{sqrt}(?discr)) * (-A - i*\text{sqrt}(?discr))$
using *assms eq-cnj-iff-real*[of A]
by *simp*
also have $\dots = (B * \text{cnj } B) * (A^2 + ?discr)$
using *complex-of-real-Re*[of A] $\langle is\text{-real } A \rangle \langle ?discr > 0 \rangle$
by (*simp add: power2-eq-square field-simps*)

finally

have $z1 * \text{cnj } z1 = (B * \text{cnj } B)^2$
using *complex-of-real-Re*[of $(\text{cmod } B)^2 - (\text{Re } A)^2$] *complex-of-real-Re*[of A] $\langle is\text{-real } A \rangle$
using *complex-mult-cnj-cmod*[of B]
by (*simp add: power2-eq-square*)

moreover

have $z2 * \text{cnj } z2 = (B * \text{cnj } B)^2$
using *assms*
by *simp*

ultimately

show *?thesis*
by *simp*

qed

lemma *calc-ideal-point1-mix*:

assumes $(z1, z2) = \text{calc-ideal-point1-cvec } A \ B \ is\text{-real } A \ (\text{cmod } B)^2 > (\text{cmod } A)^2$
shows $B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 = -2 * A * (B * \text{cnj } B)^2$

proof–

have $B * \text{cnj } z1 + \text{cnj } B * z1 = -2 * A * B * \text{cnj } B$
using *assms eq-cnj-iff-real*[of A]
by (*simp, simp add: field-simps*)

moreover

have $\text{cnj } z2 = z2$
using *assms*
by *simp*

hence $B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 = (B * \text{cnj } z1 + \text{cnj } B * z1) * z2$
by (*simp add: field-simps*)

ultimately

have $B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 = -2 * A * (B * \text{cnj } B) * z2$
by *simp*

also have $\dots = -2 * A * (B * \text{cnj } B)^2$

using *assms*

using *complex-mult-cnj-cmod*[of B]

by (*simp add: power2-eq-square*)

finally

show ?thesis

·
qed

lemma calc-ideal-point2-mix:

assumes $(z1, z2) = \text{calc-ideal-point2-cvec } A \ B \text{ is-real } A \ (\text{cmod } B)^2 > (\text{cmod } A)^2$
shows $B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 = -2 * A * (B * \text{cnj } B)^2$

proof-

have $B * \text{cnj } z1 + \text{cnj } B * z1 = -2 * A * B * \text{cnj } B$
using *assms eq-cnj-iff-real*[of *A*]
by (*simp, simp add: field-simps*)

moreover

have $\text{cnj } z2 = z2$
using *assms*
by *simp*

hence $B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 = (B * \text{cnj } z1 + \text{cnj } B * z1) * z2$
by (*simp add: field-simps*)

ultimately

have $B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 = -2 * A * (B * \text{cnj } B) * z2$
by *simp*

also have $\dots = -2 * A * (B * \text{cnj } B)^2$
using *assms*
using *complex-mult-cnj-cmod*[of *B*]
by (*simp add: power2-eq-square*)

finally

show ?thesis

·
qed

lemma calc-ideal-point1-on-circline:

assumes $(z1, z2) = \text{calc-ideal-point1-cvec } A \ B \text{ is-real } A \ (\text{cmod } B)^2 > (\text{cmod } A)^2$
shows $A * z1 * \text{cnj } z1 + B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 + A * z2 * \text{cnj } z2 = 0$ (is ?lhs = 0)

proof-

have ?lhs = $A * (z1 * \text{cnj } z1 + z2 * \text{cnj } z2) + (B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2)$
by (*simp add: field-simps*)

also have $\dots = 2 * A * (B * \text{cnj } B)^2 + (-2 * A * (B * \text{cnj } B)^2)$
using *calc-ideal-point1-sq*[OF *assms*]
using *calc-ideal-point1-mix*[OF *assms*]
by *simp*

finally

show ?thesis
by *simp*

qed

lemma calc-ideal-point2-on-circline:

assumes $(z1, z2) = \text{calc-ideal-point2-cvec } A \ B \text{ is-real } A \ (\text{cmod } B)^2 > (\text{cmod } A)^2$
shows $A * z1 * \text{cnj } z1 + B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2 + A * z2 * \text{cnj } z2 = 0$ (is ?lhs = 0)

proof-

have ?lhs = $A * (z1 * \text{cnj } z1 + z2 * \text{cnj } z2) + (B * \text{cnj } z1 * z2 + \text{cnj } B * z1 * \text{cnj } z2)$
by (*simp add: field-simps*)

also have $\dots = 2 * A * (B * \text{cnj } B)^2 + (-2 * A * (B * \text{cnj } B)^2)$
using *calc-ideal-point2-sq*[OF *assms*]
using *calc-ideal-point2-mix*[OF *assms*]
by *simp*

finally

show ?thesis
by *simp*

qed

lemma calc-ideal-points-on-circline:

assumes *is-poincare-line* *H*
shows $\forall z \in \text{calc-ideal-points } H. z \in \text{circline-set } H$
using *assms*
unfolding *circline-set-def*
apply *simp*

proof (*transfer, transfer*)

```

fix H
assume hh: hermitean H  $\wedge$  H  $\neq$  mat-zero
obtain A B C D where *: H = (A, B, C, D)
  by (cases H, auto)
obtain z11 z12 z21 z22 where **: (z11, z12) = calc-ideal-point1-cvec A B (z21, z22) = calc-ideal-point2-cvec A B
  by (cases calc-ideal-point1-cvec A B, cases calc-ideal-point2-cvec A B) auto

assume is-poincare-line-cmat H
hence  $\forall (z1, z2) \in \text{calc-ideal-points-cmat-cvec } H. A * z1 * cnj z1 + B * cnj z1 * z2 + C * z1 * cnj z2 + D * z2 * cnj z2 = 0$ 
  using * ** hh
  using hermitean-elems[of A B C D]
  using calc-ideal-point1-on-circline[of z11 z12 A B]
  using calc-ideal-point2-on-circline[of z21 z22 A B]
  by (auto simp del: calc-ideal-point1-cvec-def calc-ideal-point2-cvec-def)
thus Ball (calc-ideal-points-cmat-cvec H) (on-circline-cmat-cvec H)
  using on-circline-cmat-cvec-circline-equation *
  by (auto simp del: on-circline-cmat-cvec-def calc-ideal-points-cmat-cvec-def simp add: field-simps)
qed

```

Calculated ideal points of an h-line are different

```

lemma calc-ideal-points-cvec-different [simp]:
  assumes (cmod B)2 > (cmod A)2 is-real A
  shows  $\neg$  (calc-ideal-point1-cvec A B  $\approx_v$  calc-ideal-point2-cvec A B)
  using assms
  by (auto) (auto simp add: cmod-def)

```

```

lemma calc-ideal-points-different:
  assumes is-poincare-line H
  shows  $\exists i1 \in (\text{calc-ideal-points } H). \exists i2 \in (\text{calc-ideal-points } H). i1 \neq i2$ 
  using assms
proof (transfer, transfer)
  fix H
  assume hh: hermitean H  $\wedge$  H  $\neq$  mat-zero is-poincare-line-cmat H
  obtain A B C D where *: H = (A, B, C, D)
    by (cases H, auto)
  hence is-real A using hh hermitean-elems by auto
  thus  $\exists i1 \in \text{calc-ideal-points-cmat-cvec } H. \exists i2 \in \text{calc-ideal-points-cmat-cvec } H. \neg i1 \approx_v i2$ 
    using * hh calc-ideal-points-cvec-different[of A B]
    apply (rule-tac x=calc-ideal-point1-cvec A B in be: I)
    apply (rule-tac x=calc-ideal-point2-cvec A B in be: I)
    by auto
qed

```

```

lemma two-calc-ideal-points [simp]:
  assumes is-poincare-line H
  shows card (calc-ideal-points H) = 2
proof -
  have  $\exists x \in \text{calc-ideal-points } H. \exists y \in \text{calc-ideal-points } H. \forall z \in \text{calc-ideal-points } H. z = x \vee z = y$ 
    by (transfer, transfer, case-tac H, simp del: calc-ideal-point1-cvec-def calc-ideal-point2-cvec-def)
  then obtain x y where *: calc-ideal-points H = {x, y}
    by auto
  moreover
  have  $x \neq y$ 
    using calc-ideal-points-different[OF assms] *
    by auto
  ultimately
  show ?thesis
    by auto
qed

```

4.3.2 Ideal points

Next we give a genuine definition of ideal points – these are the intersections of the h-line with the unit circle

```

definition ideal-points :: circline  $\Rightarrow$  complex-homo set where
  ideal-points H = circline-intersection H unit-circle

```

Ideal points are on the unit circle and on the h-line

lemma *ideal-points-on-unit-circle*:
shows $\forall z \in \text{ideal-points } H. z \in \text{circline-set unit-circle}$
unfolding *ideal-points-def circline-intersection-def circline-set-def*
by *simp*

lemma *ideal-points-on-circline*:
shows $\forall z \in \text{ideal-points } H. z \in \text{circline-set } H$
unfolding *ideal-points-def circline-intersection-def circline-set-def*
by *simp*

For each h-line there are exactly two ideal points

lemma *two-ideal-points*:
assumes *is-poincare-line H*
shows $\text{card } (\text{ideal-points } H) = 2$
proof –
have $H \neq \text{unit-circle}$
using *assms not-is-poincare-line-unit-circle*
by *auto*
let $?int = \text{circline-intersection } H \text{ unit-circle}$
obtain $i1\ i2$ **where** $i1 \in ?int\ i2 \in ?int\ i1 \neq i2$
using *calc-ideal-points-on-circline[OF assms]*
using *calc-ideal-points-on-unit-circle[of H]*
using *calc-ideal-points-different[OF assms]*
unfolding *circline-intersection-def circline-set-def*
by *auto*
thus $?thesis$
unfolding *ideal-points-def*
using *circline-intersection-at-most-2-points[OF ‹H ≠ unit-circle›]*
using *card-geq-2-iff-contains-2-elems[of ?int]*
by *auto*
qed

They are exactly the two points that our calculation finds

lemma *ideal-points-unique*:
assumes *is-poincare-line H*
shows $\text{ideal-points } H = \text{calc-ideal-points } H$
proof –
have $\text{calc-ideal-points } H \subseteq \text{ideal-points } H$
using *calc-ideal-points-on-circline[OF assms]*
using *calc-ideal-points-on-unit-circle[of H]*
unfolding *ideal-points-def circline-intersection-def circline-set-def*
by *auto*
moreover
have $H \neq \text{unit-circle}$
using *not-is-poincare-line-unit-circle assms*
by *auto*
hence *finite (ideal-points H)*
using *circline-intersection-at-most-2-points[of H unit-circle]*
unfolding *ideal-points-def*
by *auto*
ultimately
show $?thesis$
using *card-subset-eq[of ideal-points H calc-ideal-points H]*
using *two-calc-ideal-points[OF assms]*
using *two-ideal-points[OF assms]*
by *auto*
qed

For each h-line we can obtain two different ideal points

lemma *obtain-ideal-points*:
assumes *is-poincare-line H*
obtains $i1\ i2$ **where** $i1 \neq i2\ \text{ideal-points } H = \{i1, i2\}$
using *two-ideal-points[OF assms] card-eq-2-iff-doubleton[of ideal-points H]*
by *blast*

Ideal points of each h-line constructed from two points in the disc are different than those two points

lemma *ideal-points-different*:
assumes $u \in \text{unit-disc } v \in \text{unit-disc } u \neq v$
assumes *ideal-points* (*poincare-line* $u v$) = $\{i1, i2\}$
shows $i1 \neq i2 \ u \neq i1 \ u \neq i2 \ v \neq i1 \ v \neq i2$
proof –
have $i1 \in \text{ocircline-set } \text{ounit-circle } i2 \in \text{ocircline-set } \text{ounit-circle}$
using *assms(3) assms(4) ideal-points-on-unit-circle is-poincare-line-poincare-line*
by *fastforce+*
thus $u \neq i1 \ u \neq i2 \ v \neq i1 \ v \neq i2$
using *assms(1–2)*
using *disc-inter-ocircline-set[of ounit-circle]*
unfolding *unit-disc-def*
by *auto*
show $i1 \neq i2$
using *assms*
by (*metis doubleton-eq-iff is-poincare-line-poincare-line obtain-ideal-points*)
qed

H-line is uniquely determined by its ideal points

lemma *ideal-points-line-unique*:
assumes *is-poincare-line* H *ideal-points* $H = \{i1, i2\}$
shows $H = \text{poincare-line } i1 \ i2$
by (*smt assms(1) assms(2) calc-ideal-points-on-unit-circle circline-set-def ex-poincare-line-points ideal-points-different(1) ideal-points-on-circline ideal-points-unique insertI1 insert-commute inversion-unit-circle mem-Collect-eq unique-poincare-line-general*)

Ideal points of some special h-lines

Ideal points of *x-axis*

lemma *ideal-points-x-axis*
[*simp*]: *ideal-points* *x-axis* = $\{\text{of-complex } (-1), \text{of-complex } 1\}$
proof (*subst ideal-points-unique, simp*)
have *calc-ideal-points-clmat-hcoords* *x-axis-clmat* = $\{\text{of-complex-hcoords } (-1), \text{of-complex-hcoords } 1\}$
by *transfer auto*
thus *calc-ideal-points* *x-axis* = $\{\text{of-complex } (-1), \text{of-complex } 1\}$
by (*simp add: calc-ideal-points.abs-eq of-complex.abs-eq x-axis-def*)
qed

Ideal points are proportional vectors only if h-line is a line segment passing through zero

lemma *ideal-points-proportional*:
assumes *is-poincare-line* H *ideal-points* $H = \{i1, i2\}$ *to-complex* $i1 = \text{cor } k * \text{to-complex } i2$
shows $0_h \in \text{circline-set } H$
proof –
have $i1 \neq i2$
using $\langle \text{ideal-points } H = \{i1, i2\} \rangle$
using $\langle \text{is-poincare-line } H \rangle$ *ex-poincare-line-points ideal-points-different(1)* **by** *blast*

have $i1 \in \text{circline-set } \text{unit-circle } i2 \in \text{circline-set } \text{unit-circle}$
using *assms calc-ideal-points-on-unit-circle ideal-points-unique*
by *blast+*

hence *cmod* (*cor* k) = 1
using $\langle \text{to-complex } i1 = \text{cor } k * \text{to-complex } i2 \rangle$
by (*metis (mono-tags, lifting) circline-set-unit-circle imageE mem-Collect-eq mult.right-neutral norm-mult to-complex-of-complex unit-circle-set-def*)
hence $k = -1$
using $\langle \text{to-complex } i1 = \text{cor } k * \text{to-complex } i2 \rangle \langle i1 \neq i2 \rangle$
using $\langle i1 \in \text{circline-set } \text{unit-circle} \rangle \langle i2 \in \text{circline-set } \text{unit-circle} \rangle$
by (*smt (verit, best) mult-cancel-right1 norm-of-real not-inf-on-unit-circle'' of-complex-to-complex of-real-1*)

have $\forall i1 \in \text{calc-ideal-points } H. \forall i2 \in \text{calc-ideal-points } H. \text{is-poincare-line } H \wedge i1 \neq i2 \wedge \text{to-complex } i1 = - \text{to-complex } i2 \longrightarrow$
 $0_h \in \text{circline-set } H$
unfolding *circline-set-def*
proof (*simp, transfer, transfer, safe*)

```

fix A B C D i11 i12 i21 i22 k
assume H:hermitean (A, B, C, D) (A, B, C, D) ≠ mat-zero
assume line: is-poincare-line-cmat (A, B, C, D)
assume i1: (i11, i12) ∈ calc-ideal-points-cmat-cvec (A, B, C, D)
assume i2:(i21, i22) ∈ calc-ideal-points-cmat-cvec (A, B, C, D)
assume ¬ (i11, i12) ≈v (i21, i22)
assume opposite: to-complex-cvec (i11, i12) = - to-complex-cvec (i21, i22)

```

```

let ?discr = sqrt ((cmod B)2 - (Re D)2)
let ?den = (cmod B)2
let ?i1 = B * (- D - i * ?discr)
let ?i2 = B * (- D + i * ?discr)

```

```

have i11 = ?i1 ∨ i11 = ?i2 i12 = ?den
  i21 = ?i1 ∨ i21 = ?i2 i22 = ?den
  using i1 i2 H line
  by (auto split: if-split-asm)
hence i: i11 = ?i1 ∧ i21 = ?i2 ∨ i11 = ?i2 ∧ i21 = ?i1
  using ⟨¬ (i11, i12) ≈v (i21, i22)⟩
  by auto

```

```

have ?den ≠ 0
  using line
  by auto

```

```

hence i11 = - i21
  using opposite ⟨i12 = ?den⟩ ⟨i22 = ?den⟩
  by (simp add: nonzero-neg-divide-eq-eq2)

```

```

hence ?i1 = - ?i2
  using i
  by (metis add.inverse-inverse)

```

```

hence D = 0
  using ⟨?den ≠ 0⟩
  by (simp add: field-simps)

```

```

thus on-circline-cmat-cvec (A, B, C, D) 0v
  by (simp add: vec-cnj-def)

```

qed

```

thus ?thesis
  using assms ⟨k = -1⟩
  using calc-ideal-points-different ideal-points-unique
  by fastforce

```

qed

Transformations of ideal points

Möbius transformations that fix the unit disc when acting on h-lines map their ideal points to ideal points.

lemma *ideal-points-moebius-circline* [simp]:

assumes *unit-circle-fix* M *is-poincare-line* H

shows *ideal-points* (moebius-circline M H) = (moebius-pt M) ‘ (ideal-points H) (is ?I' = ?M ‘ ?I)

proof—

obtain i1 i2 **where** *: i1 ≠ i2 ?I = {i1, i2}

using assms(2)

by (rule obtain-ideal-points)

let ?Mi1 = ?M i1 **and** ?Mi2 = ?M i2

have ?Mi1 ∈ ?M ‘ (circline-set H)

?Mi2 ∈ ?M ‘ (circline-set H)

?Mi1 ∈ ?M ‘ (circline-set unit-circle)

?Mi2 ∈ ?M ‘ (circline-set unit-circle)

using *

unfolding *ideal-points-def* *circline-intersection-def* *circline-set-def*

by blast+

hence $?Mi1 \in ?I'$
 $?Mi2 \in ?I'$
using *unit-circle-fix-iff*[of M] *assms*
unfolding *ideal-points-def* *circline-intersection-def* *circline-set-def*
by (*metis mem-Collect-eq* *moebius-circline*)+
moreover
have $?Mi1 \neq ?Mi2$
using *bij-moebius-pt*[of M] *
using *moebius-pt-invert* **by** *blast*
moreover
have *is-poincare-line* (*moebius-circline* M H)
using *assms* *unit-circle-fix-preserve-is-poincare-line*
by *simp*
ultimately
have $?I' = \{?Mi1, ?Mi2\}$
using *two-ideal-points*[of *moebius-circline* M H]
using *card-eq-2-doubleton*[of $?I'$ $?Mi1$ $?Mi2$]
by *simp*
thus *?thesis*
using *(2)
by *auto*
qed

lemma *ideal-points-poincare-line-moebius* [*simp*]:
assumes *unit-disc-fix* M $u \in \text{unit-disc}$ $v \in \text{unit-disc}$ $u \neq v$
assumes *ideal-points* (*poincare-line* u v) = $\{i1, i2\}$
shows *ideal-points* (*poincare-line* (*moebius-pt* M u) (*moebius-pt* M v)) = $\{\text{moebius-pt } M \ i1, \text{moebius-pt } M \ i2\}$
using *assms*
by *auto*

Conjugation also maps ideal points to ideal points

lemma *ideal-points-conjugate* [*simp*]:
assumes *is-poincare-line* H
shows *ideal-points* (*conjugate-circline* H) = *conjugate* ‘ (*ideal-points* H) (**is** $?I' = ?M$ ‘ $?I$)
proof–
obtain $i1$ $i2$ **where** *: $i1 \neq i2$ $?I = \{i1, i2\}$
using *assms*
by (*rule obtain-ideal-points*)
let $?Mi1 = ?M \ i1$ **and** $?Mi2 = ?M \ i2$
have $?Mi1 \in ?M$ ‘ (*circline-set* H)
 $?Mi2 \in ?M$ ‘ (*circline-set* H)
 $?Mi1 \in ?M$ ‘ (*circline-set* *unit-circle*)
 $?Mi2 \in ?M$ ‘ (*circline-set* *unit-circle*)
using *
unfolding *ideal-points-def* *circline-intersection-def* *circline-set-def*
by *blast*+
hence $?Mi1 \in ?I'$
 $?Mi2 \in ?I'$
unfolding *ideal-points-def* *circline-intersection-def* *circline-set-def*
using *circline-set-conjugate-circline* *circline-set-def* *conjugate-unit-circle-set*
by *blast*+
moreover
have $?Mi1 \neq ?Mi2$
using $\langle i1 \neq i2 \rangle$
by (*auto simp add: conjugate-inj*)
moreover
have *is-poincare-line* (*conjugate-circline* H)
using *assms*
by *simp*
ultimately
have $?I' = \{?Mi1, ?Mi2\}$
using *two-ideal-points*[of *conjugate-circline* H]
using *card-eq-2-doubleton*[of $?I'$ $?Mi1$ $?Mi2$]
by *simp*
thus *?thesis*
using *(2)

by auto
qed

lemma *ideal-points-poincare-line-conjugate* [simp]:
assumes $u \in \text{unit-disc } v \in \text{unit-disc } u \neq v$
assumes *ideal-points* (poincare-line $u v$) = { $i1, i2$ }
shows *ideal-points* (poincare-line (conjugate u) (conjugate v)) = {conjugate $i1$, conjugate $i2$ }
using *assms*
by auto

end
theory *Poincare-Distance*
imports *Poincare-Lines-Ideal-Points Hyperbolic-Functions*
begin

5 H-distance in the Poincaré model

Informally, the *h-distance* between the two h-points is defined as the absolute value of the logarithm of the cross ratio between those two points and the two ideal points.

abbreviation *Re-cross-ratio* **where** *Re-cross-ratio* $z u v w \equiv \text{Re } (\text{to-complex } (\text{cross-ratio } z u v w))$

definition *calc-poincare-distance* :: *complex-homo* \Rightarrow *complex-homo* \Rightarrow *complex-homo* \Rightarrow *complex-homo* \Rightarrow *real* **where**
[*simp*]: *calc-poincare-distance* $u i1 v i2 = \text{abs } (\ln (\text{Re-cross-ratio } u i1 v i2))$

definition *poincare-distance-pred* :: *complex-homo* \Rightarrow *complex-homo* \Rightarrow *real* \Rightarrow *bool* **where**
[*simp*]: *poincare-distance-pred* $u v d \longleftrightarrow$
 $(u = v \wedge d = 0) \vee (u \neq v \wedge (\forall i1 i2. \text{ideal-points } (\text{poincare-line } u v) = \{i1, i2\} \longrightarrow d = \text{calc-poincare-distance } u i1 v i2))$

definition *poincare-distance* :: *complex-homo* \Rightarrow *complex-homo* \Rightarrow *real* **where**
poincare-distance $u v = (\text{THE } d. \text{poincare-distance-pred } u v d)$

We shown that the described cross-ratio is always finite, positive real number.

lemma *distance-cross-ratio-real-positive*:
assumes $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$ **and** $u \neq v$
shows $\forall i1 i2. \text{ideal-points } (\text{poincare-line } u v) = \{i1, i2\} \longrightarrow$
 $\text{cross-ratio } u i1 v i2 \neq \infty_h \wedge \text{is-real } (\text{to-complex } (\text{cross-ratio } u i1 v i2)) \wedge \text{Re-cross-ratio } u i1 v i2 > 0$

(is ?P $u v$)

proof (rule *wlog-positive-x-axis*[OF *assms*])

fix x

assume *: *is-real* x $0 < \text{Re } x$ $\text{Re } x < 1$

hence $x \neq -1$ $x \neq 1$

by auto

hence **: *of-complex* $x \neq \infty_h$ *of-complex* $x \neq 0_h$ *of-complex* $x \neq \text{of-complex } (-1)$ *of-complex* $1 \neq \text{of-complex } x$
of-complex $x \in \text{circline-set } x\text{-axis}$

using *

unfolding *circline-set-x-axis*

by (auto *simp add: of-complex-inj*)

have ***: $0_h \neq \text{of-complex } (-1)$ $0_h \neq \text{of-complex } 1$
by (*metis of-complex-zero-iff zero-neq-neg-one, simp*)

have ****: $-x - 1 \neq 0$ $x - 1 \neq 0$
using $\langle x \neq -1 \rangle \langle x \neq 1 \rangle$
by (*metis add.inverse-inverse eq-iff-diff-eq-0, simp*)

have *poincare-line* 0_h (*of-complex* x) = *x-axis*

using **

by (*simp add: poincare-line-0-real-is-x-axis*)

thus ?P 0_h (*of-complex* x)

using * ** **** *****

using *cross-ratio-not-inf*[*of* 0_h *of-complex* 1 *of-complex* (-1) *of-complex* x]

using *cross-ratio-not-inf*[*of* 0_h *of-complex* (-1) *of-complex* 1 *of-complex* x]

using *cross-ratio-real*[*of* 0 -1 x 1] *cross-ratio-real*[*of* 0 1 x -1]

```

apply (auto simp add: poincare-line-0-real-is-x-axis doubleton-eq-iff circline-set-x-axis)
apply (subst cross-ratio, simp-all, subst Re-complex-div-gt-0, simp, subst mult-neg-neg, simp-all)+
done
next
fix M u v
let ?Mu = moebius-pt M u and ?Mv = moebius-pt M v
assume *: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc u ≠ v
        ?P ?Mu ?Mv
show ?P u v
proof safe
  fix i1 i2
  let ?cr = cross-ratio u i1 v i2
  assume **: ideal-points (poincare-line u v) = {i1, i2}
  have i1 ≠ u i1 ≠ v i2 ≠ u i2 ≠ v i1 ≠ i2
    using ideal-points-different[OF *(2-3), of i1 i2] ** ⟨u ≠ v⟩
    by auto
  hence 0 < Re (to-complex ?cr) ∧ is-real (to-complex ?cr) ∧ ?cr ≠ ∞h
    using * **
  apply (erule-tac x=moebius-pt M i1 in allE)
  apply (erule-tac x=moebius-pt M i2 in allE)
  apply (subst (asm) ideal-points-poincare-line-moebius[of M u v i1 i2], simp-all)
  done
  thus 0 < Re (to-complex ?cr) is-real (to-complex ?cr) ?cr = ∞h ⇒ False
    by simp-all
qed
qed

```

Next we can show that for every different points from the unit disc there is exactly one number that satisfies the h-distance predicate.

lemma *distance-unique*:

```

assumes u ∈ unit-disc and v ∈ unit-disc
shows ∃! d. poincare-distance-pred u v d
proof (cases u = v)
  case True
  thus ?thesis
    by auto
next
  case False
  obtain i1 i2 where *: i1 ≠ i2 ideal-points (poincare-line u v) = {i1, i2}
    using obtain-ideal-points[OF is-poincare-line-poincare-line] ⟨u ≠ v⟩
    by blast
  let ?d = calc-poincare-distance u i1 v i2
  show ?thesis
  proof (rule ex1I)
    show poincare-distance-pred u v ?d
      using * ⟨u ≠ v⟩
    proof (simp del: calc-poincare-distance-def, safe)
      fix i1' i2'
      assume {i1, i2} = {i1', i2'}
      hence **: (i1' = i1 ∧ i2' = i2) ∨ (i1' = i2 ∧ i2' = i1)
        using doubleton-eq-iff[of i1 i2 i1' i2']
        by blast
      have all-different: u ≠ i1 u ≠ i2 v ≠ i1 v ≠ i2 u ≠ i1' u ≠ i2' v ≠ i1' v ≠ i2' i1 ≠ i2
        using ideal-points-different[OF assms, of i1 i2] * ** ⟨u ≠ v⟩
        by auto

      show calc-poincare-distance u i1 v i2 = calc-poincare-distance u i1' v i2'
    proof–
      let ?cr = cross-ratio u i1 v i2
      let ?cr' = cross-ratio u i1' v i2'

      have Re (to-complex ?cr) > 0 is-real (to-complex ?cr)
        Re (to-complex ?cr') > 0 is-real (to-complex ?cr')
        using False distance-cross-ratio-real-positive[OF assms(1-2)] * **
        by auto
    
```

```

thus ?thesis
  using **
  using cross-ratio-not-zero cross-ratio-not-inf all-different
  by auto (subst cross-ratio-commute-24, subst reciprocal-real, simp-all add: ln-div)
qed
qed
next
fix d
assume poincare-distance-pred u v d
thus d = ?d
  using * ⟨u ≠ v⟩
  by auto
qed
qed

```

```

lemma poincare-distance-satisfies-pred [simp]:
assumes u ∈ unit-disc and v ∈ unit-disc
shows poincare-distance-pred u v (poincare-distance u v)
  using distance-unique[OF assms] theI'[of poincare-distance-pred u v]
  unfolding poincare-distance-def
  by blast

```

```

lemma poincare-distance-I:
assumes u ∈ unit-disc and v ∈ unit-disc and u ≠ v and ideal-points (poincare-line u v) = {i1, i2}
shows poincare-distance u v = calc-poincare-distance u i1 v i2
using assms
using poincare-distance-satisfies-pred[OF assms(1-2)]
by simp

```

```

lemma poincare-distance-refl [simp]:
assumes u ∈ unit-disc
shows poincare-distance u u = 0
using assms
using poincare-distance-satisfies-pred[OF assms assms]
by simp

```

Unit disc preserving Möbius transformations preserve h-distance.

```

lemma unit-disc-fix-preserve-poincare-distance [simp]:
assumes unit-disc-fix M and u ∈ unit-disc and v ∈ unit-disc
shows poincare-distance (moebius-pt M u) (moebius-pt M v) = poincare-distance u v
proof (cases u = v)
case True
have moebius-pt M u ∈ unit-disc moebius-pt M v ∈ unit-disc
  using unit-disc-fix-iff[OF assms(1), symmetric] assms
  by blast+
thus ?thesis
  using assms ⟨u = v⟩
  by simp
next
case False
obtain i1 i2 where *: ideal-points (poincare-line u v) = {i1, i2}
  using ⟨u ≠ v⟩
  by (rule obtain-ideal-points[OF is-poincare-line-poincare-line[of u v]])
let ?Mu = moebius-pt M u and ?Mv = moebius-pt M v and ?Mi1 = moebius-pt M i1 and ?Mi2 = moebius-pt M i2

have **: ?Mu ∈ unit-disc ?Mv ∈ unit-disc
  using assms
  using unit-disc-fix-iff
  by blast+

have ***: ?Mu ≠ ?Mv
  using ⟨u ≠ v⟩
  by simp

have poincare-distance u v = calc-poincare-distance u i1 v i2
  using poincare-distance-I[OF assms(2-3) ⟨u ≠ v⟩ *]

```

```

  by auto
moreover
have unit-circle-fix M
  using assms
  by simp
hence ++: ideal-points (poincare-line ?Mu ?Mv) = {?Mi1, ?Mi2}
  using ⟨u ≠ v⟩ assms *
  by simp
have poincare-distance ?Mu ?Mv = calc-poincare-distance ?Mu ?Mi1 ?Mv ?Mi2
  by (rule poincare-distance-I[OF ** *** ++])
moreover
have calc-poincare-distance ?Mu ?Mi1 ?Mv ?Mi2 = calc-poincare-distance u i1 v i2
  using ideal-points-different[OF assms(2-3) ⟨u ≠ v⟩ *]
  unfolding calc-poincare-distance-def
  by (subst moebius-preserve-cross-ratio[symmetric], simp-all)
ultimately
show ?thesis
  by simp
qed

```

Knowing ideal points for x-axis, we can easily explicitly calculate distances.

lemma *poincare-distance-x-axis-x-axis*:

assumes $x \in \text{unit-disc}$ **and** $y \in \text{unit-disc}$ **and** $x \in \text{circline-set } x\text{-axis}$ **and** $y \in \text{circline-set } x\text{-axis}$
shows $\text{poincare-distance } x \ y =$

(let $x' = \text{to-complex } x$; $y' = \text{to-complex } y$
 in $\text{abs} (\ln (\text{Re} (((1 + x') * (1 - y')) / ((1 - x') * (1 + y'))))))$)

proof–

obtain $x' \ y'$ **where** $*$: $x = \text{of-complex } x'$ $y = \text{of-complex } y'$
 using $\text{inf-or-of-complex}[of \ x] \ \text{inf-or-of-complex}[of \ y] \ \langle x \in \text{unit-disc} \rangle \ \langle y \in \text{unit-disc} \rangle$
 by auto

have $\text{cmod } x' < 1 \ \text{cmod } y' < 1$
 using $\langle x \in \text{unit-disc} \rangle \ \langle y \in \text{unit-disc} \rangle *$
 by (metis *unit-disc-iff-cmod-lt-1*)
hence $**$: $x' \neq 1 \ x' \neq -1 \ y' \neq 1 \ y' \neq -1$
 by auto

have $1 + y' \neq 0$
 using $**$
 by (metis *add.left-cancel add-neg-numeral-special*(7))

show *?thesis*

proof (*cases* $x = y$)

case *True*

thus *?thesis*

using *assms*(1-2)
 using *unit-disc-iff-cmod-lt-1*[*of to-complex* x] $*$ $**$ $\langle 1 + y' \neq 0 \rangle$
 by auto

next

case *False*

hence *poincare-line* $x \ y = x\text{-axis}$

using *poincare-line-x-axis*[*OF* *assms*]

by *simp*

hence *ideal-points* (*poincare-line* $x \ y$) = {*of-complex* (-1) , *of-complex* 1 }

by *simp*

hence *poincare-distance* $x \ y = \text{calc-poincare-distance } x \ (\text{of-complex } (-1)) \ y \ (\text{of-complex } 1)$

using *poincare-distance-I* *assms* $\langle x \neq y \rangle$

by *auto*

also have $\dots = \text{abs} (\ln (\text{Re} (((x' + 1) * (y' - 1)) / ((x' - 1) * (y' + 1))))))$

using $*$ $\langle \text{cmod } x' < 1 \rangle \ \langle \text{cmod } y' < 1 \rangle$

by (*simp*, *transfer*, *transfer*, *auto*)

finally

show *?thesis*

using $*$

by (metis (*no-types*, *lifting*) *add commute minus-diff-eq minus-divide-divide mult-minus-left mult-minus-right*)

to-complex-of-complex)

qed
qed

lemma *poincare-distance-zero-x-axis*:

assumes $x \in \text{unit-disc}$ **and** $x \in \text{circline-set } x\text{-axis}$

shows $\text{poincare-distance } 0_h x = (\text{let } x' = \text{to-complex } x \text{ in } \text{abs } (\ln (\text{Re } ((1 - x') / (1 + x')))))$

using *assms*

using *poincare-distance-x-axis-x-axis*[of $0_h x$]

by (*simp add: Let-def*)

lemma *poincare-distance-zero*:

assumes $x \in \text{unit-disc}$

shows $\text{poincare-distance } 0_h x = (\text{let } x' = \text{to-complex } x \text{ in } \text{abs } (\ln (\text{Re } ((1 - \text{cmod } x') / (1 + \text{cmod } x')))))$ (**is** $?P x$)

proof (*cases* $x = 0_h$)

case *True*

thus $?thesis$

by *auto*

next

case *False*

show $?thesis$

proof (*rule wlog-rotation-to-positive-x-axis*)

show $x \in \text{unit-disc } x \neq 0_h$ **by** *fact+*

next

fix φu

assume $u \in \text{unit-disc } u \neq 0_h$ $?P (\text{moebius-pt } (\text{moebius-rotation } \varphi) u)$

thus $?P u$

using *unit-disc-fix-preserve-poincare-distance*[of *moebius-rotation* φ $0_h u$]

by (*cases* $u = \infty_h$) (*simp-all add: Let-def*)

next

fix x

assume *is-real* x $0 < \text{Re } x$ $\text{Re } x < 1$

thus $?P (\text{of-complex } x)$

using *poincare-distance-zero-x-axis*[of *of-complex* x]

by *simp (auto simp add: circline-set-x-axis cmod-eq-Re complex-is-Real-iff)*

qed

qed

lemma *poincare-distance-zero-opposite* [*simp*]:

assumes *of-complex* $z \in \text{unit-disc}$

shows $\text{poincare-distance } 0_h (\text{of-complex } (-z)) = \text{poincare-distance } 0_h (\text{of-complex } z)$

proof–

have $*$: *of-complex* $(-z) \in \text{unit-disc}$

using *assms*

by *auto*

show $?thesis$

using *poincare-distance-zero*[*OF* *assms*]

using *poincare-distance-zero*[*OF* $*$]

by *simp*

qed

5.1 Distance explicit formula

Instead of the h-distance itself, very frequently its hyperbolic cosine is analyzed.

abbreviation *cosh-dist* $u v \equiv \text{cosh } (\text{poincare-distance } u v)$

lemma *cosh-poincare-distance-cross-ratio-average*:

assumes $u \in \text{unit-disc } v \in \text{unit-disc } u \neq v$ *ideal-points* $(\text{poincare-line } u v) = \{i1, i2\}$

shows $\text{cosh-dist } u v =$

$$((\text{Re-cross-ratio } u i1 v i2) + (\text{Re-cross-ratio } v i1 u i2)) / 2$$

proof–

let $?cr = \text{cross-ratio } u i1 v i2$

let $?crRe = \text{Re } (\text{to-complex } ?cr)$

have $?cr \neq \infty_h$ *is-real* $(\text{to-complex } ?cr)$ $?crRe > 0$

using *distance-cross-ratio-real-positive*[*OF* *assms*(1–3)] *assms*(4)

by *simp-all*

then obtain cr **where** $*$: *cross-ratio* $u\ i1\ v\ i2 = of\ complex\ cr\ cr \neq 0\ is\ real\ cr\ Re\ cr > 0$
using *inf-or-of-complex*[of *cross-ratio* $u\ i1\ v\ i2$]
by (*smt to-complex-of-complex zero-complex.simps*(1))
thus *?thesis*
using $*$
using *assms cross-ratio-commute-13*[of $v\ i1\ u\ i2$]
unfolding *poincare-distance-I*[*OF assms*] *calc-poincare-distance-def cosh-def*
by (*cases* $Re\ cr \geq 1$)
(auto simp add: ln-div[of 0] *exp-minus field-simps Re-divide power2-eq-square complex.expand)*
qed

definition *poincare-distance-formula'* :: *complex* \Rightarrow *complex* \Rightarrow *real* **where**
*[simp]: poincare-distance-formula' u v = 1 + 2 * ((cmod (u - v))² / ((1 - (cmod u)²) * (1 - (cmod v)²)))*

Next we show that the following formula expresses h-distance between any two h-points (note that the ideal points do not figure anymore).

definition *poincare-distance-formula* :: *complex* \Rightarrow *complex* \Rightarrow *real* **where**
[simp]: poincare-distance-formula u v = arcosh (poincare-distance-formula' u v)

lemma *blaschke-preserve-distance-formula* [*simp*]:

assumes *of-complex* $k \in unit-disc\ u \in unit-disc\ v \in unit-disc$

shows *poincare-distance-formula (to-complex (moebius-pt (blaschke k) u)) (to-complex (moebius-pt (blaschke k) v)) = poincare-distance-formula (to-complex u) (to-complex v)*

proof (*cases* $k = 0$)

case *True*

thus *?thesis*

by *simp*

next

case *False*

obtain $u'\ v'$ **where** $*$: $u' = to-complex\ u\ v' = to-complex\ v$

by *auto*

have $cmod\ u' < 1\ cmod\ v' < 1\ cmod\ k < 1$

using *assms* $*$

using *inf-or-of-complex*[of u] *inf-or-of-complex*[of v]

by *auto*

obtain $nu\ du\ nv\ dv\ d\ kk\ ddu\ ddv$ **where**

$**$: $nu = u' - k\ du = 1 - cnj\ k * u'\ nv = v' - k\ dv = 1 - cnj\ k * v'$

$d = u' - v'\ ddu = 1 - u' * cnj\ u'\ ddv = 1 - v' * cnj\ v'\ kk = 1 - k * cnj\ k$

by *auto*

have $d: nu * dv - nv * du = d * kk$

by (*subst* $**$) + (*simp add: field-simps*)

have $ddu: du * cnj\ du - nu * cnj\ nu = ddu * kk$

by (*subst* $**$) + (*simp add: field-simps*)

have $ddv: dv * cnj\ dv - nv * cnj\ nv = ddv * kk$

by (*subst* $**$) + (*simp add: field-simps*)

have $du \neq 0$

proof (*rule ccontr*)

assume $\neg ?thesis$

hence $cmod (1 - cnj\ k * u') = 0$

using $\langle du = 1 - cnj\ k * u' \rangle$

by *auto*

hence $cmod (cnj\ k * u') = 1$

by *auto*

thus *False*

using $\langle cmod\ k < 1 \rangle\ \langle cmod\ u' < 1 \rangle$

using *mult-strict-mono*[of $cmod\ k\ 1\ cmod\ u'\ 1$]

by (*simp add: norm-mult*)

qed

have $dv \neq 0$

proof (*rule ccontr*)

assume $\neg ?thesis$

hence $cmod (1 - cnj k * v') = 0$
using $\langle dv = 1 - cnj k * v' \rangle$
by *auto*
hence $cmod (cnj k * v') = 1$
by *auto*
thus *False*
using $\langle cmod k < 1 \rangle \langle cmod v' < 1 \rangle$
using *mult-strict-mono*[of $cmod k 1 cmod v' 1$]
by (*simp add: norm-mult*)
qed

have $kk \neq 0$
proof (*rule ccontr*)
assume $\neg ?thesis$
hence $cmod (1 - k * cnj k) = 0$
using $\langle kk = 1 - k * cnj k \rangle$
by *auto*
hence $cmod (k * cnj k) = 1$
by *auto*
thus *False*
using $\langle cmod k < 1 \rangle$
using *mult-strict-mono*[of $cmod k 1 cmod k 1$]
using *complex-mod-sqrt-Re-mult-cnj* **by** *auto*
qed

note $nz = \langle du \neq 0 \rangle \langle dv \neq 0 \rangle \langle kk \neq 0 \rangle$

have $nu / du - nv / dv = (nu*dv - nv*du) / (du * dv)$
using *nz*
by (*simp add: field-simps*)
hence $(cmod (nu/du - nv/dv))^2 = cmod ((d*kk) / (du*dv) * (cnj ((d*kk) / (du*dv))))$ (*is ?lhs = -*)
unfolding *complex-mod-mult-cnj*[*symmetric*]
by (*subst (asm) d*) *simp*
also have $\dots = cmod ((d*cnj d*kk*kk) / (du*cnj du*dv*cnj dv))$
by (*simp add: field-simps norm-mult norm-divide*)
finally have $1: ?lhs = cmod ((d*cnj d*kk*kk) / (du*cnj du*dv*cnj dv))$.

have $(1 - ((cmod nu) / (cmod du))^2) * (1 - ((cmod nv) / (cmod dv))^2) =$
 $(1 - cmod((nu * cnj nu) / (du * cnj du))) * (1 - cmod((nv * cnj nv) / (dv * cnj dv)))$ (*is ?rhs = -*)
by (*metis norm-divide complex-mod-mult-cnj power-divide*)
also have $\dots = cmod(((du*cnj du - nu*cnj nu) / (du * cnj du)) * ((dv*cnj dv - nv*cnj nv) / (dv * cnj dv)))$
proof-

have $u' \neq 1 / cnj k \vee v' \neq 1 / cnj k$
using $\langle cmod u' < 1 \rangle \langle cmod v' < 1 \rangle \langle cmod k < 1 \rangle$
by (*auto simp add: False norm-divide*)

moreover

have $cmod k \neq 1$
using $\langle cmod k < 1 \rangle$
by *linarith*

ultimately

have $cmod (nu/du) < 1 \wedge cmod (nv/dv) < 1$

using $** (1-4)$

using *unit-disc-fix-discI*[*OF blaschke-unit-disc-fix*[*OF* $\langle cmod k < 1 \rangle$] $\langle u \in unit-disc \rangle$] $\langle u' = to-complex u \rangle$

using *unit-disc-fix-discI*[*OF blaschke-unit-disc-fix*[*OF* $\langle cmod k < 1 \rangle$] $\langle v \in unit-disc \rangle$] $\langle v' = to-complex v \rangle$

using *inf-or-of-complex*[of $u \in unit-disc$] *inf-or-of-complex*[of $v \in unit-disc$]

using *moebius-pt-blaschke*[of $k u$] **using** *moebius-pt-blaschke*[of $k v$]

by *auto*

hence $(cmod (nu/du))^2 < 1 \wedge (cmod (nv/dv))^2 < 1$

by (*simp-all add: cmod-def*)

hence $cmod (nu * cnj nu / (du * cnj du)) < 1 \wedge cmod (nv * cnj nv / (dv * cnj dv)) < 1$

by (*metis complex-mod-mult-cnj norm-divide power-divide*)**+**

moreover

have *is-real* $(nu * cnj nu / (du * cnj du)) \wedge$ *is-real* $(nv * cnj nv / (dv * cnj dv))$

using *eq-cnj-iff-real*[of $nu * cnj nu / (du * cnj du)$]

using *eq-cnj-iff-real*[of $nv * cnj nv / (dv * cnj dv)$]

```

  by (auto simp add: mult.commute)
moreover
have Re (nu * cnj nu / (du * cnj du)) ≥ 0 Re (nv * cnj nv / (dv * cnj dv)) ≥ 0
  using ⟨du ≠ 0⟩ ⟨dv ≠ 0⟩
  unfolding complex-mult-cnj-cmod
  by simp-all
ultimately
have 1 - cmod (nu * cnj nu / (du * cnj du)) = cmod (1 - nu * cnj nu / (du * cnj du))
  1 - cmod (nv * cnj nv / (dv * cnj dv)) = cmod (1 - nv * cnj nv / (dv * cnj dv))
  by (simp-all add: cmod-def)
thus ?thesis
  using nz
  by (simp add: diff-divide-distrib norm-mult)
qed
also have ... = cmod(((ddu * kk) / (du * cnj du)) * ((ddv * kk) / (dv * cnj dv)))
  by (subst ddu, subst ddv, simp)
also have ... = cmod((ddu*ddv*kk*kk) / (du*cnj du*dv*cnj dv))
  by (simp add: field-simps)
finally have 2: ?rhs = cmod((ddu*ddv*kk*kk) / (du*cnj du*dv*cnj dv))
.

have ?lhs / ?rhs =
  cmod ((d*cnj d*kk*kk) / (du*cnj du*dv*cnj dv)) / cmod((ddu*ddv*kk*kk) / (du*cnj du*dv*cnj dv))
  by (subst 1, subst 2, simp)
also have ... = cmod ((d*cnj d)/(ddu*ddv))
  using nz by (simp add: norm-mult norm-divide)
also have ... = (cmod d)2 / ((1 - (cmod u)2)*(1 - (cmod v)2))
proof-
  have (cmod u)2 < 1 (cmod v)2 < 1
    using ⟨cmod u' < 1⟩ ⟨cmod v' < 1⟩
    by (simp-all add: cmod-def)
  hence cmod (1 - u' * cnj u') = 1 - (cmod u')2 cmod (1 - v' * cnj v') = 1 - (cmod v')2
    by (auto simp add: cmod-eq-Re cmod-power2 power2-eq-square[symmetric])
  thus ?thesis
    using nz
    by (simp add: *(6) *(7) norm-divide norm-mult power2-eq-square)
qed
finally
have 3: ?lhs / ?rhs = (cmod d)2 / ((1 - (cmod u)2)*(1 - (cmod v)2)) .

have cmod k ≠ 1 u' ≠ 1 / cnj k v' ≠ 1 / cnj k u ≠ ∞h v ≠ ∞h
  using ⟨cmod k < 1⟩ ⟨u ∈ unit-disc⟩ ⟨v ∈ unit-disc⟩ * ⟨k ≠ 0⟩ ** ⟨kk ≠ 0⟩ nz
  by auto
thus ?thesis using assms
  using * ** 3
  using moebius-pt-blaschke[of k u]
  using moebius-pt-blaschke[of k v']
  by (simp add: norm-divide)
qed

```

To prove the equivalence between the h-distance definition and the distance formula, we shall employ the without loss of generality principle. Therefore, we must show that the distance formula is preserved by h-isometries.

Rotation preserve *poincare-distance-formula*.

lemma *rotation-preserve-distance-formula* [simp]:

```

  assumes u ∈ unit-disc v ∈ unit-disc
  shows poincare-distance-formula (to-complex (moebius-pt (moebius-rotation φ) u)) (to-complex (moebius-pt (moebius-rotation φ) v)) =
    poincare-distance-formula (to-complex u) (to-complex v)
  using assms
  using inf-or-of-complex[of u] inf-or-of-complex[of v]
  by (auto simp: norm-mult)

```

Unit disc fixing Möbius preserve *poincare-distance-formula*.

lemma *unit-disc-fix-preserve-distance-formula* [simp]:

```

  assumes unit-disc-fix M u ∈ unit-disc v ∈ unit-disc

```

shows *poincare-distance-formula* (*to-complex* (*moebius-pt* M u)) (*to-complex* (*moebius-pt* M v)) =
poincare-distance-formula (*to-complex* u) (*to-complex* v) (is ? P' u v M)

proof–
have $\forall u \in \text{unit-disc. } \forall v \in \text{unit-disc. } ?P' u v M$ (is ? P M)
proof (*rule wlog-unit-disc-fix*[*OF assms*(1)])
fix k
assume *cmod* $k < 1$
hence *of-complex* $k \in \text{unit-disc}$
by *simp*
thus ? P (*blaschke* k)
using *blaschke-preserve-distance-formula*
by *simp*

next
fix φ
show ? P (*moebius-rotation* φ)
using *rotation-preserve-distance-formula*
by *simp*

next
fix $M1$ $M2$
assume *: ? P $M1$ **and** **: ? P $M2$ **and** *u11*: *unit-disc-fix* $M1$ *unit-disc-fix* $M2$
thus ? P ($M1 + M2$)
by (*auto simp del: poincare-distance-formula-def*)

qed
thus ?*thesis*
using *assms*
by *simp*

qed

The equivalence between the two h-distance representations.

lemma *poincare-distance-formula*:

assumes $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$
shows *poincare-distance* u v = *poincare-distance-formula* (*to-complex* u) (*to-complex* v) (is ? P u v)

proof (*rule wlog-x-axis*)
fix x
assume *: *is-real* x $0 \leq \text{Re } x$ $\text{Re } x < 1$
show ? P 0_h (*of-complex* x) (is ?*lhs* = ?*rhs*)

proof–
have *of-complex* $x \in \text{unit-disc}$ *of-complex* $x \in \text{circline-set } x\text{-axis}$ *cmod* $x < 1$
using * *cmod-eq-Re*
by (*auto simp add: circline-set-x-axis*)
hence ?*lhs* = $|\ln (\text{Re} ((1 - x) / (1 + x)))|$
using *poincare-distance-zero-x-axis*[*of of-complex* x]
by *simp*

moreover
have ?*rhs* = $|\ln (\text{Re} ((1 - x) / (1 + x)))|$

proof–
let ? $x = 1 + 2 * (\text{cmod } x)^2 / (1 - (\text{cmod } x)^2)$
have $0 \leq 2 * (\text{cmod } x)^2 / (1 - (\text{cmod } x)^2)$
by (*smt* $\langle \text{cmod } x < 1 \rangle$ *divide-nonneg-nonneg norm-ge-zero power-le-one zero-le-power2*)
hence *arcosh-real-gt*: $1 \leq ?x$
by *auto*
have ?*rhs* = *arcosh* ? x
by *simp*

also have ... = $\ln ((1 + (\text{cmod } x)^2) / (1 - (\text{cmod } x)^2) + 2 * (\text{cmod } x) / (1 - (\text{cmod } x)^2))$

proof–
have $1 - (\text{cmod } x)^2 > 0$
using $\langle \text{cmod } x < 1 \rangle$
by (*smt norm-not-less-zero one-power2 power2-eq-imp-eq power-mono*)
hence 1: ? $x = (1 + (\text{cmod } x)^2) / (1 - (\text{cmod } x)^2)$
by (*simp add: field-simps*)
have 2: ? $x^2 - 1 = (4 * (\text{cmod } x)^2) / (1 - (\text{cmod } x)^2)^2$
using $\langle 1 - (\text{cmod } x)^2 > 0 \rangle$
apply (*subst* 1)
unfolding *power-divide*
by (*subst divide-diff-eq-iff, simp, simp add: power2-eq-square field-simps*)
show ?*thesis*

```

    using ⟨1 - (cmod x)2 > 0⟩
    apply (subst arcosh-real-def[OF arcosh-real-gt])
    apply (subst 2)
    apply (subst 1)
    apply (subst real-sqrt-divide)
    apply (subst real-sqrt-mult)
    apply simp
  done
qed
also have ... = ln (((1 + (cmod x))2) / (1 - (cmod x)2))
  apply (subst add-divide-distrib[symmetric])
  apply (simp add: field-simps power2-eq-square)
  done
also have ... = ln ((1 + cmod x) / (1 - (cmod x)))
  using ⟨cmod x < 1⟩
  using square-diff-square-factored[of 1 cmod x]
  by (simp add: power2-eq-square)
also have ... = |ln (Re ((1 - x) / (1 + x)))|
proof-
  have *: Re ((1 - x) / (1 + x)) ≤ 1 Re ((1 - x) / (1 + x)) > 0
    using ⟨is-real x⟩ ⟨Re x ≥ 0⟩ ⟨Re x < 1⟩
    using complex-is-Real-iff
    by auto
  hence |ln (Re ((1 - x) / (1 + x)))| = - ln (Re ((1 - x) / (1 + x)))
    by auto
  hence |ln (Re ((1 - x) / (1 + x)))| = ln (Re ((1 + x) / (1 - x)))
    using ln-div[of 1 Re ((1 - x)/(1 + x))] * ⟨is-real x⟩
    by (fastforce simp: complex-is-Real-iff)
  moreover
  have ln ((1 + cmod x) / (1 - cmod x)) = ln ((1 + Re x) / (1 - Re x))
    using ⟨Re x ≥ 0⟩ ⟨is-real x⟩
    using cmod-eq-Re by auto
  moreover
  have (1 + Re x) / (1 - Re x) = Re ((1 + x) / (1 - x))
    using ⟨is-real x⟩ ⟨Re x < 1⟩
    by (smt Re-divide-real eq-iff-diff-eq-0 minus-complex.simps one-complex.simps plus-complex.simps)
  ultimately
  show ?thesis
    by simp
qed
finally
show ?thesis
  .
qed
ultimately
show ?thesis
  by simp
qed
next
fix M u v
assume *: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc
assume ?P (moebius-pt M u) (moebius-pt M v)
thus ?P u v
  using *(1-3)
  by (simp del: poincare-distance-formula-def)
next
show u ∈ unit-disc v ∈ unit-disc
  by fact+
qed

```

Some additional properties proved easily using the distance formula.

poincare-distance is symmetric.

lemma *poincare-distance-sym*:

assumes $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$

shows $\text{poincare-distance } u \ v = \text{poincare-distance } v \ u$

```

using assms
using poincare-distance-formula[OF assms(1) assms(2)]
using poincare-distance-formula[OF assms(2) assms(1)]
by (simp add: mult.commute norm.minus.commute)

```

```

lemma poincare-distance-formula'-ge-1:
assumes  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
shows  $1 \leq \text{poincare-distance-formula}'(\text{to-complex } u) (\text{to-complex } v)$ 
using unit-disc-cmod-square-lt-1[OF assms(1)] unit-disc-cmod-square-lt-1[OF assms(2)]
by auto

```

poincare-distance is non-negative.

```

lemma poincare-distance-ge0:
assumes  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
shows  $\text{poincare-distance } u \ v \geq 0$ 
using poincare-distance-formula'-ge-1 assms by (simp add: poincare-distance-formula)

```

```

lemma cosh-dist:
assumes  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
shows  $\text{cosh-dist } u \ v = \text{poincare-distance-formula}'(\text{to-complex } u) (\text{to-complex } v)$ 
using poincare-distance-formula[OF assms] poincare-distance-formula'-ge-1[OF assms]
by simp

```

poincare-distance is zero only if the two points are equal.

```

lemma poincare-distance-eq-0-iff:
assumes  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
shows  $\text{poincare-distance } u \ v = 0 \iff u = v$ 
using assms
apply auto
using poincare-distance-formula'-ge-1[OF assms]
using unit-disc-cmod-square-lt-1[OF assms(1)] unit-disc-cmod-square-lt-1[OF assms(2)]
apply (simp add: poincare-distance-formula)
by (simp add: unit-disc-to-complex-inj)

```

Conjugate preserve *poincare-distance-formula*.

```

lemma conjugate-preserve-poincare-distance [simp]:
assumes  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
shows  $\text{poincare-distance } (\text{conjugate } u) (\text{conjugate } v) = \text{poincare-distance } u \ v$ 
proof–
obtain  $u' \ v'$  where  $u = \text{of-complex } u'$   $v = \text{of-complex } v'$ 
using assms inf-or-of-complex[of u] inf-or-of-complex[of v]
by auto

```

```

have **:  $\text{conjugate } u \in \text{unit-disc}$   $\text{conjugate } v \in \text{unit-disc}$ 
using * assms
by auto

```

```

show ?thesis
using *
using poincare-distance-formula[OF assms]
using poincare-distance-formula[OF **]

```

```

by (metis complex-cnj-diff complex-mod-cnj conjugate-of-complex poincare-distance-def poincare-distance-formula'-def
poincare-distance-formula-def to-complex-of-complex)
qed

```

5.2 Existence and uniqueness of points with a given distance

```

lemma ex-x-axis-poincare-distance-negative':
fixes  $d :: \text{real}$ 
assumes  $d \geq 0$ 
shows  $\text{let } z = (1 - \exp d) / (1 + \exp d)$ 
in  $\text{is-real } z \wedge \text{Re } z \leq 0 \wedge \text{Re } z > -1 \wedge$ 
 $\text{of-complex } z \in \text{unit-disc} \wedge \text{of-complex } z \in \text{circline-set } x\text{-axis} \wedge$ 
 $\text{poincare-distance } 0_{\text{h}} (\text{of-complex } z) = d$ 
proof–

```

```

have exp d ≥ 1
  using assms
  using one-le-exp-iff[of d, symmetric]
  by blast

hence 1 + exp d ≠ 0
  by linarith

let ?z = (1 - exp d) / (1 + exp d)

have ?z ≤ 0
  using ⟨exp d ≥ 1⟩
  by (simp add: divide-nonpos-nonneg)

moreover

have ?z > -1
  using exp-gt-zero[of d]
  by (smt divide-less-eq-1-neg nonzero-minus-divide-right)

moreover

hence abs ?z < 1
  using ⟨?z ≤ 0⟩
  by simp
hence cmod ?z < 1
  by (metis norm-of-real)
hence of-complex ?z ∈ unit-disc
  by simp

moreover
have of-complex ?z ∈ circline-set x-axis
  unfolding circline-set-x-axis
  by simp

moreover
have (1 - ?z) / (1 + ?z) = exp d
proof-
  have 1 + ?z = 2 / (1 + exp d)
    using ⟨1 + exp d ≠ 0⟩
    by (subst add-divide-eq-iff, auto)
  moreover
  have 1 - ?z = 2 * exp d / (1 + exp d)
    using ⟨1 + exp d ≠ 0⟩
    by (subst diff-divide-eq-iff, auto)
  ultimately
  show ?thesis
    using ⟨1 + exp d ≠ 0⟩
    by simp
qed

ultimately
show ?thesis
  using poincare-distance-zero-x-axis[of of-complex ?z]
  using ⟨d ≥ 0⟩ ⟨exp d ≥ 1⟩
  by simp (simp add: cmod-eq-Re)
qed

lemma ex-x-axis-poincare-distance-negative:
  assumes d ≥ 0
  shows ∃ z. is-real z ∧ Re z ≤ 0 ∧ Re z > -1 ∧
    of-complex z ∈ unit-disc ∧ of-complex z ∈ circline-set x-axis ∧
    poincare-distance 0_h (of-complex z) = d (is ∃ z. ?P z)
  using ex-x-axis-poincare-distance-negative'[OF assms]
  unfolding Let-def
  by blast

```

For each real number d there is exactly one point on the positive x-axis such that h-distance between 0 and that point is d .

lemma *unique-x-axis-poincare-distance-negative*:

assumes $d \geq 0$

shows $\exists! z. \text{is-real } z \wedge \text{Re } z \leq 0 \wedge \text{Re } z > -1 \wedge$
 $\text{poincare-distance } 0_h (\text{of-complex } z) = d$ (**is** $\exists! z. ?P z$)

proof–

let $?z = (1 - \exp d) / (1 + \exp d)$

have $?P ?z$

using *ex-x-axis-poincare-distance-negative* [OF assms]

unfolding *Let-def*

by *blast*

moreover

have $\forall z'. ?P z' \longrightarrow z' = ?z$

proof–

let $?g = \lambda x'. |\ln (\text{Re } ((1 - x') / (1 + x')))|$

let $?A = \{x. \text{is-real } x \wedge \text{Re } x > -1 \wedge \text{Re } x \leq 0\}$

have *inj-on* (*poincare-distance* $0_h \circ \text{of-complex}$) $?A$

proof (*rule comp-inj-on*)

show *inj-on of-complex* $?A$

using *of-complex-inj*

unfolding *inj-on-def*

by *blast*

next

show *inj-on* (*poincare-distance* 0_h) (*of-complex* $?A$) (**is** *inj-on* $?f$ (*of-complex* $?A$))

proof (*subst inj-on-cong*)

have $*$: *of-complex* $?A =$

$\{z. z \in \text{unit-disc} \wedge z \in \text{circline-set } x\text{-axis} \wedge \text{Re } (\text{to-complex } z) \leq 0\}$ (**is** $- = ?B$)

by (*auto simp add: cmod-eq-Re circline-set-x-axis*)

fix x

assume $x \in \text{of-complex } ?A$

hence $x \in ?B$

using $*$

by *simp*

thus *poincare-distance* $0_h x = (?g \circ \text{to-complex}) x$

using *poincare-distance-zero-x-axis*

by (*simp add: Let-def*)

next

have $*$: *to-complex* $?A = \text{of-complex } ?A$

by (*auto simp add: image-iff*)

show *inj-on* ($?g \circ \text{to-complex}$) (*of-complex* $?A$)

proof (*rule comp-inj-on*)

show *inj-on to-complex* (*of-complex* $?A$)

unfolding *inj-on-def*

by *auto*

next

have *inj-on* $?g ?A$

unfolding *inj-on-def*

proof(*safe*)

fix $x y$

assume hh : *is-real* x *is-real* $y - 1 < \text{Re } x \text{Re } x \leq 0$

$- 1 < \text{Re } y \text{Re } y \leq 0 \quad |\ln (\text{Re } ((1 - x) / (1 + x)))| = |\ln (\text{Re } ((1 - y) / (1 + y)))|$

have *is-real* $((1 - x)/(1 + x))$

using $\langle \text{is-real } x \rangle \text{div-reals}[of 1-x 1+x]$

by *auto*

have *is-real* $((1 - y)/(1 + y))$

using $\langle \text{is-real } y \rangle \text{div-reals}[of 1-y 1+y]$

by *auto*

have $\text{Re } (1 + x) > 0$


```

    using ⟨- 1 < Re x⟩ by auto
  hence 1 + x ≠ 0
    by force
  have Re (1 - x) ≥ 0
    using ⟨Re x ≤ 0⟩ by auto
  hence Re ((1 - x)/(1 + x)) > 0
    using Re-divide-real ⟨0 < Re (1 + x)⟩ complex-eq-if-Re-eq hh(1) hh(4) by auto
  have Re(1 - x) ≥ Re (1 + x)
    using hh by auto
  hence Re ((1 - x)/(1 + x)) ≥ 1
    using ⟨Re (1 + x) > 0⟩ ⟨is-real ((1 - x)/(1 + x))⟩
  by (smt Re-divide-real arg-0-iff hh(1) le-divide-eq-1-pos one-complex.simps(2) plus-complex.simps(2))

  have Re (1 + y) > 0
    using ⟨- 1 < Re y⟩ by auto
  hence 1 + y ≠ 0
    by force
  have Re (1 - y) ≥ 0
    using ⟨Re y ≤ 0⟩ by auto
  hence Re ((1 - y)/(1 + y)) > 0
    using Re-divide-real ⟨0 < Re (1 + y)⟩ complex-eq-if-Re-eq hh by auto
  have Re(1 - y) ≥ Re (1 + y)
    using hh by auto
  hence Re ((1 - y)/(1 + y)) ≥ 1
    using ⟨Re (1 + y) > 0⟩ ⟨is-real ((1 - y)/(1 + y))⟩
  by (smt Re-divide-real arg-0-iff hh le-divide-eq-1-pos one-complex.simps(2) plus-complex.simps(2))

  have ln (Re ((1 - x) / (1 + x))) = ln (Re ((1 - y) / (1 + y)))
    using ⟨Re ((1 - y)/(1 + y)) ≥ 1⟩ ⟨Re ((1 - x)/(1 + x)) ≥ 1⟩ hh
    by auto
  hence Re ((1 - x) / (1 + x)) = Re ((1 - y) / (1 + y))
    using ⟨Re ((1 - y)/(1 + y)) > 0⟩ ⟨Re ((1 - x)/(1 + x)) > 0⟩
    by auto
  hence (1 - x) / (1 + x) = (1 - y) / (1 + y)
    using ⟨is-real ((1 - y)/(1 + y))⟩ ⟨is-real ((1 - x)/(1 + x))⟩
    using complex-eq-if-Re-eq by blast
  hence (1 - x) * (1 + y) = (1 - y) * (1 + x)
    using ⟨1 + y ≠ 0⟩ ⟨1 + x ≠ 0⟩
    by (simp add:field-simps)
  thus x = y
    by (simp add:field-simps)
qed
thus inj-on ?g (to-complex ‘ of-complex ‘ ?A)
  using *
  by simp
qed
qed
thus ?thesis
  using ⟨?P ?z⟩
  unfolding inj-on-def
  by auto
qed
ultimately
show ?thesis
  by blast
qed

lemma ex-x-axis-poincare-distance-positive:
  assumes d ≥ 0
  shows ∃ z. is-real z ∧ Re z ≥ 0 ∧ Re z < 1 ∧
    of-complex z ∈ unit-disc ∧ of-complex z ∈ circline-set x-axis ∧
    poincare-distance 0h (of-complex z) = d (is ∃ z. is-real z ∧ Re z ≥ 0 ∧ Re z < 1 ∧ ?P z)
proof-
  obtain z where *: is-real z Re z ≤ 0 Re z > -1 ?P z
    using ex-x-axis-poincare-distance-negative[OF assms]

```

```

  by auto
  hence **: of-complex  $z \in \text{unit-disc}$  of-complex  $z \in \text{circline-set } x\text{-axis}$ 
    by (auto simp add: cmod-eq-Re)
  have is-real  $(-z) \wedge \text{Re } (-z) \geq 0 \wedge \text{Re } (-z) < 1 \wedge ?P (-z)$ 
    using * **
    by (simp add: circline-set-x-axis)
  thus ?thesis
    by blast
qed

```

lemma *unique-x-axis-poincare-distance-positive:*

```

  assumes  $d \geq 0$ 
  shows  $\exists! z. \text{is-real } z \wedge \text{Re } z \geq 0 \wedge \text{Re } z < 1 \wedge$ 
     $\text{poincare-distance } 0_h (\text{of-complex } z) = d$  (is  $\exists! z. \text{is-real } z \wedge \text{Re } z \geq 0 \wedge \text{Re } z < 1 \wedge ?P z$ )

```

proof–

```

  obtain  $z$  where *: is-real  $z$   $\text{Re } z \leq 0$   $\text{Re } z > -1$   $?P z$ 
    using unique-x-axis-poincare-distance-negative[OF assms]
    by auto
  hence **: of-complex  $z \in \text{unit-disc}$  of-complex  $z \in \text{circline-set } x\text{-axis}$ 
    by (auto simp add: cmod-eq-Re circline-set-x-axis)
  show ?thesis

```

proof

```

  show is-real  $(-z) \wedge \text{Re } (-z) \geq 0 \wedge \text{Re } (-z) < 1 \wedge ?P (-z)$ 
    using * **
    by simp

```

next

```

  fix  $z'$ 
  assume is-real  $z' \wedge \text{Re } z' \geq 0 \wedge \text{Re } z' < 1 \wedge ?P z'$ 
  hence is-real  $(-z') \wedge \text{Re } (-z') \leq 0 \wedge \text{Re } (-z') > -1 \wedge ?P (-z')$ 
    by (auto simp add: circline-set-x-axis cmod-eq-Re)
  hence  $-z' = z$ 
    using unique-x-axis-poincare-distance-negative[OF assms] *
    by blast
  thus  $z' = -z$ 
    by auto

```

qed

qed

Equal distance implies that segments are isometric - this means that congruence could be defined either by two segments having the same distance or by requiring existence of an isometry that maps one segment to the other.

lemma *poincare-distance-eq-ex-moebius:*

```

  assumes in-disc:  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$  and  $u' \in \text{unit-disc}$  and  $v' \in \text{unit-disc}$ 
  assumes poincare-distance  $u v = \text{poincare-distance } u' v'$ 
  shows  $\exists M. \text{unit-disc-fix } M \wedge \text{moebius-pt } M u = u' \wedge \text{moebius-pt } M v = v'$  (is  $?P' u v u' v'$ )

```

proof (cases $u = v$)

case *True*

thus ?thesis

```

  using assms poincare-distance-eq-0-iff[of  $u' v'$ ]
  by (simp add: unit-disc-fix-transitive)

```

next

case *False*

```

  have  $\forall u' v'. u \neq v \wedge u' \in \text{unit-disc} \wedge v' \in \text{unit-disc} \wedge \text{poincare-distance } u v = \text{poincare-distance } u' v' \longrightarrow$ 
     $?P' u' v' u v$  (is  $?P u v$ )

```

proof (rule wlog-positive-x-axis[where $P = ?P$])

fix x

assume is-real x $0 < \text{Re } x$ $\text{Re } x < 1$

hence of-complex $x \in \text{unit-disc}$ of-complex $x \in \text{circline-set } x\text{-axis}$

unfolding *circline-set-x-axis*

by (auto simp add: cmod-eq-Re)

show $?P 0_h (\text{of-complex } x)$

proof *safe*

fix $u' v'$

assume $0_h \neq \text{of-complex } x$ and in-disc: $u' \in \text{unit-disc}$ $v' \in \text{unit-disc}$ and

$\text{poincare-distance } 0_h (\text{of-complex } x) = \text{poincare-distance } u' v'$

hence $u' \neq v'$ $\text{poincare-distance } u' v' > 0$

```

using poincare-distance-eq-0-iff[of  $0_h$  of-complex  $x$ ] ⟨of-complex  $x \in \text{unit-disc}$ ⟩
using poincare-distance-ge0[of  $0_h$  of-complex  $x$ ]
by auto
then obtain  $M$  where  $M$ : unit-disc-fix  $M$  moebius-pt  $M$   $u' = 0_h$  moebius-pt  $M$   $v' \in \text{positive-x-axis}$ 
using ex-unit-disc-fix-to-zero-positive-x-axis[of  $u' v'$ ] in-disc
by auto

then obtain  $Mv'$  where  $Mv'$ : moebius-pt  $M$   $v' = \text{of-complex } Mv'$ 
using inf-or-of-complex[of moebius-pt  $M$   $v'$ ] in-disc unit-disc-fix-iff[of  $M$ ]
by (metis image-eqI inf-notin-unit-disc)

have moebius-pt  $M$   $v' \in \text{unit-disc}$ 
using  $M(1)$  ⟨ $v' \in \text{unit-disc}$ ⟩
by auto

have  $\text{Re } Mv' > 0$  is-real  $Mv'$   $\text{Re } Mv' < 1$ 
using  $M$   $Mv'$  of-complex-inj ⟨moebius-pt  $M$   $v' \in \text{unit-disc}$ ⟩
unfolding positive-x-axis-def circline-set-x-axis
using cmod-eq-Re
by auto fastforce

have poincare-distance  $0_h$  (moebius-pt  $M$   $v'$ ) = poincare-distance  $u' v'$ 
using  $M(1)$ 
using in-disc
by (subst  $M(2)$ )[symmetric], simp

have  $Mv' = x$ 
using ⟨poincare-distance  $0_h$  (moebius-pt  $M$   $v'$ ) = poincare-distance  $u' v'$ ⟩  $Mv'$ 
using ⟨poincare-distance  $0_h$  (of-complex  $x$ ) = poincare-distance  $u' v'$ ⟩
using unique-x-axis-poincare-distance-positive[of poincare-distance  $u' v'$ ]
⟨poincare-distance  $u' v' > 0$ ⟩
using ⟨ $\text{Re } Mv' > 0$ ⟩ ⟨ $\text{Re } Mv' < 1$ ⟩ ⟨is-real  $Mv'$ ⟩
using ⟨is-real  $x$ ⟩ ⟨ $\text{Re } x > 0$ ⟩ ⟨ $\text{Re } x < 1$ ⟩
unfolding positive-x-axis-def
by auto

thus  $?P' u' v' 0_h$  (of-complex  $x$ )
using  $M$   $Mv'$ 
by auto
qed
next
show  $u \in \text{unit-disc}$   $v \in \text{unit-disc}$   $u \neq v$ 
by fact+
next
fix  $M$   $u$   $v$ 
let  $?Mu = \text{moebius-pt } M$   $u$  and  $?Mv = \text{moebius-pt } M$   $v$ 
assume 1: unit-disc-fix  $M$   $u \in \text{unit-disc}$   $v \in \text{unit-disc}$   $u \neq v$ 
hence 2:  $?Mu \neq ?Mv$   $?Mu \in \text{unit-disc}$   $?Mv \in \text{unit-disc}$ 
by auto
assume 3:  $?P$  (moebius-pt  $M$   $u$ ) (moebius-pt  $M$   $v$ )
show  $?P$   $u$   $v$ 
proof safe
fix  $u' v'$ 
assume 4:  $u' \in \text{unit-disc}$   $v' \in \text{unit-disc}$  poincare-distance  $u$   $v = \text{poincare-distance } u' v'$ 
hence poincare-distance  $?Mu$   $?Mv = \text{poincare-distance } u$   $v$ 
using 1
by simp
then obtain  $M'$  where 5: unit-disc-fix  $M'$  moebius-pt  $M'$   $u' = ?Mu$  moebius-pt  $M'$   $v' = ?Mv$ 
using 2 3 4
by auto
let  $?M = (-M) + M'$ 
have unit-disc-fix  $?M \wedge \text{moebius-pt } ?M$   $u' = u \wedge \text{moebius-pt } ?M$   $v' = v$ 
using 5 ⟨unit-disc-fix  $M$ ⟩
using unit-disc-fix-moebius-comp[of  $-M$   $M'$ ]
using unit-disc-fix-moebius-inv[of  $M$ ]
by simp

```

thus $\exists M. \text{unit-disc-fix } M \wedge \text{moebius-pt } M u' = u \wedge \text{moebius-pt } M v' = v$
by *blast*
qed
qed
then obtain M **where** $\text{unit-disc-fix } M \wedge \text{moebius-pt } M u' = u \wedge \text{moebius-pt } M v' = v$
using *assms* $\langle u \neq v \rangle$
by *blast*
hence $\text{unit-disc-fix } (-M) \wedge \text{moebius-pt } (-M) u = u' \wedge \text{moebius-pt } (-M) v = v'$
using *unit-disc-fix-moebius-inv*[*of* M]
by *auto*
thus *?thesis*
by *blast*
qed

lemma *unique-midpoint-x-axis*:
assumes x : *is-real* $x - 1 < \text{Re } x \text{ Re } x < 1$ **and**
 y : *is-real* $y - 1 < \text{Re } y \text{ Re } y < 1$ **and**
 $x \neq y$
shows $\exists! z. -1 < \text{Re } z \wedge \text{Re } z < 1 \wedge \text{is-real } z \wedge \text{poincare-distance (of-complex } z) (\text{of-complex } x) = \text{poincare-distance (of-complex } z) (\text{of-complex } y)$ **(is** $\exists! z. ?R z (\text{of-complex } x) (\text{of-complex } y)$ **)**

proof–
let $?x = \text{of-complex } x$ **and** $?y = \text{of-complex } y$
let $?P = \lambda x y. \exists! z. ?R z x y$
have $\forall x. -1 < \text{Re } x \wedge \text{Re } x < 1 \wedge \text{is-real } x \wedge \text{of-complex } x \neq ?y \longrightarrow ?P (\text{of-complex } x) ?y$ **(is** $?Q (\text{of-complex } y)$ **)**
proof (*rule wlog-real-zero*)
show $?y \in \text{unit-disc}$
using y
by (*simp add: cmod-eq-Re*)
next
show *is-real* (*to-complex* $?y$)
using y
by *simp*
next
show $?Q 0_h$
proof (*rule allI, rule impI, (erule conjE)+*)
fix x
assume $x: -1 < \text{Re } x \text{ Re } x < 1$ *is-real* x
let $?x = \text{of-complex } x$
assume $?x \neq 0_h$
hence $x \neq 0$
by *auto*
hence $\text{Re } x \neq 0$
using x
using *complex-neq-0*
by *auto*

have $*$: $\forall a. -1 < a \wedge a < 1 \longrightarrow$
 $(\text{poincare-distance (of-complex (cor } a)) ?x = \text{poincare-distance (of-complex (cor } a)) 0_h \iff$
 $(\text{Re } x) * a * a - 2 * a + \text{Re } x = 0)$

proof (*rule allI, rule impI*)
fix $a :: \text{real}$
assume $-1 < a \wedge a < 1$
hence *of-complex* (*cor* a) $\in \text{unit-disc}$
by *auto*
moreover
have $(a - \text{Re } x)^2 / ((1 - a^2) * (1 - (\text{Re } x)^2)) = a^2 / (1 - a^2) \iff$
 $(\text{Re } x) * a * a - 2 * a + \text{Re } x = 0$ **(is** $?lhs \iff ?rhs$ **)**

proof–
have $1 - a^2 \neq 0$
using $\langle -1 < a \wedge a < 1 \rangle$
by (*metis cancel-comm-monoid-add-class.diff-cancel diff-eq-diff-less less-numeral-extra(4) power2-eq-1-iff right-minus-eq*)
hence $?lhs \iff (a - \text{Re } x)^2 / (1 - (\text{Re } x)^2) = a^2$
by (*smt divide-cancel-right divide-divide-eq-left mult.commute*)
also have $\dots \iff (a - \text{Re } x)^2 = a^2 * (1 - (\text{Re } x)^2)$
proof–

```

have 1 - (Re x)2 ≠ 0
  using x
  by (smt power2-eq-1-iff)
thus ?thesis
  by (simp add: divide-eq-eq)
qed
also have ... ↔ a2 * (Re x)2 - 2*a*Re x + (Re x)2 = 0
  by (simp add: power2-diff field-simps)
also have ... ↔ Re x * (a2 * Re x - 2 * a + Re x) = 0
  by (simp add: power2-eq-square field-simps)
also have ... ↔ ?rhs
  using ⟨Re x ≠ 0⟩
  by (simp add: mult.commute mult.left-commute power2-eq-square)
finally
show ?thesis
.
qed
moreover
have arcosh (1 + 2 * ((a - Re x)2 / ((1 - a2) * (1 - (Re x)2))) = arcosh (1 + 2 * a2 / (1 - a2)) ↔ ?lhs
  using ⟨-1 < a ∧ a < 1⟩ x mult-left-cancel[of 2::real (a - Re x)2 / ((1 - a2) * (1 - (Re x)2)) a2 / (1 - a2)]
  by (subst arcosh-eq-iff, simp-all add: square-le-1)
ultimately
show poincare-distance (of-complex (cor a)) (of-complex x) = poincare-distance (of-complex (cor a)) 0h ↔
  (Re x) * a * a - 2 * a + Re x = 0
  using x
  by (auto simp add: poincare-distance-formula cmod-eq-Re)
qed

show ?P ?x 0h
proof
let ?a = (1 - sqrt(1 - (Re x)2)) / (Re x)
let ?b = (1 + sqrt(1 - (Re x)2)) / (Re x)

have is-real ?a
  by simp
moreover
have 1 - (Re x)2 > 0
  using x
  by (smt power2-eq-1-iff square-le-1)
have |?a| < 1
proof (cases Re x > 0)
case True
have (1 - Re x)2 < 1 - (Re x)2
  using ⟨Re x > 0⟩ x
  by (simp add: power2-eq-square field-simps)
hence 1 - Re x < sqrt (1 - (Re x)2)
  using real-less-rsqrt by fastforce
thus ?thesis
  using ⟨1 - (Re x)2 > 0⟩ ⟨Re x > 0⟩
  by simp
next
case False
hence Re x < 0
  using ⟨Re x ≠ 0⟩
  by simp

have 1 + Re x > 0
  using ⟨Re x > -1⟩
  by simp
hence 2*Re x + 2*Re x*Re x < 0
  using ⟨Re x < 0⟩
  by (metis comm-semiring-class.distrib mult.commute mult-2-right mult-less-0-iff one-add-one zero-less-double-add-iff-zero-less-si)
hence (1 + Re x)2 < 1 - (Re x)2
  by (simp add: power2-eq-square field-simps)
hence 1 + Re x < sqrt (1 - (Re x)2)
  using ⟨1 - (Re x)2 > 0⟩

```

```

    using real-less-rsqr by blast
  thus ?thesis
    using ⟨Re x < 0⟩
    by (simp add: field-simps)
qed
hence  $-1 < ?a \ \ ?a < 1$ 
  by linarith+
moreover
have  $(Re\ x) * ?a * ?a - 2 * ?a + Re\ x = 0$ 
  using ⟨Re x ≠ 0⟩ ⟨ $1 - (Re\ x)^2 > 0$ ⟩
  by (simp add: field-simps power2-eq-square)
ultimately
show  $-1 < Re\ (cor\ ?a) \wedge Re\ (cor\ ?a) < 1 \wedge is-real\ ?a \wedge poincare-distance\ (of-complex\ ?a)\ (of-complex\ x) =$ 
  poincare-distance (of-complex ?a) 0h
  using *
  by auto

fix z
assume **:  $-1 < Re\ z \wedge Re\ z < 1 \wedge is-real\ z \wedge$ 
  poincare-distance (of-complex z) (of-complex x) = poincare-distance (of-complex z) 0h
hence  $Re\ x * Re\ z * Re\ z - 2 * Re\ z + Re\ x = 0$ 
  using *[rule-format, of Re z] x
  by auto
moreover
have  $\sqrt{4 - 4 * Re\ x * Re\ x} = 2 * \sqrt{1 - Re\ x * Re\ x}$ 
proof-
  have  $\sqrt{4 - 4 * Re\ x * Re\ x} = \sqrt{4 * (1 - Re\ x * Re\ x)}$ 
  by simp
  thus ?thesis
  by (simp only: real-sqrt-mult, simp)
qed
moreover
have  $(2 - 2 * \sqrt{1 - Re\ x * Re\ x}) / (2 * Re\ x) = ?a$ 
proof-
  have  $(2 - 2 * \sqrt{1 - Re\ x * Re\ x}) / (2 * Re\ x) =$ 
     $(2 * (1 - \sqrt{1 - Re\ x * Re\ x})) / (2 * Re\ x)$ 
  by simp
  thus ?thesis
  by (subst (asm) mult-divide-mult-cancel-left) (auto simp add: power2-eq-square)
qed
moreover
have  $(2 + 2 * \sqrt{1 - Re\ x * Re\ x}) / (2 * Re\ x) = ?b$ 
proof-
  have  $(2 + 2 * \sqrt{1 - Re\ x * Re\ x}) / (2 * Re\ x) =$ 
     $(2 * (1 + \sqrt{1 - Re\ x * Re\ x})) / (2 * Re\ x)$ 
  by simp
  thus ?thesis
  by (subst (asm) mult-divide-mult-cancel-left) (auto simp add: power2-eq-square)
qed
ultimately
have  $Re\ z = ?a \vee Re\ z = ?b$ 
  using discriminant-nonneg[of Re x - 2 Re x Re z] discrim-def[of Re x - 2 Re x]
  using ⟨Re x ≠ 0⟩ ⟨ $-1 < Re\ x$ ⟩ ⟨ $Re\ x < 1$ ⟩ ⟨ $1 - (Re\ x)^2 > 0$ ⟩
  by (auto simp add: power2-eq-square)
have  $|?b| > 1$ 
proof (cases Re x > 0)
  case True
  have  $(Re\ x - 1)^2 < 1 - (Re\ x)^2$ 
  using ⟨Re x > 0⟩ x
  by (simp add: power2-eq-square field-simps)
  hence  $Re\ x - 1 < \sqrt{1 - (Re\ x)^2}$ 
  using real-less-rsqr
  by simp
  thus ?thesis
  using ⟨ $1 - (Re\ x)^2 > 0$ ⟩ ⟨Re x > 0⟩
  by simp

```

```

next
  case False
  hence  $Re\ x < 0$ 
    using  $\langle Re\ x \neq 0 \rangle$ 
    by simp
  have  $1 + Re\ x > 0$ 
    using  $\langle Re\ x > -1 \rangle$ 
    by simp
  hence  $2 * Re\ x + 2 * Re\ x * Re\ x < 0$ 
    using  $\langle Re\ x < 0 \rangle$ 
  by (metis comm-semiring-class.distrib mult.commute mult-2-right mult-less-0-iff one-add-one zero-less-double-add-iff-zero-less-s)
  hence  $1 - (Re\ x)^2 > (-1 - (Re\ x))^2$ 
    by (simp add: field-simps power2-eq-square)
  hence  $\sqrt{1 - (Re\ x)^2} > -1 - Re\ x$ 
    using real-less-rsqrt
    by simp
  thus ?thesis
    using  $\langle Re\ x < 0 \rangle$ 
    by (simp add: field-simps)
qed
hence  $?b < -1 \vee ?b > 1$ 
  by auto

hence  $Re\ z = ?a$ 
  using  $\langle Re\ z = ?a \vee Re\ z = ?b \rangle **$ 
  by auto
thus  $z = ?a$ 
  using ** complex-of-real-Re
  by fastforce
qed
qed
next
fix  $a\ u$ 
let  $?M = moebius-pt\ (blaschke\ a)$ 
let  $?Mu = ?M\ u$ 
assume  $u \in unit-disc\ is-real\ a\ cmod\ a < 1$ 
assume  $*$ :  $?Q\ ?Mu$ 
show  $?Q\ u$ 
proof (rule allI, rule impI, (erule conjE)+)
  fix  $x$ 
  assume  $x: -1 < Re\ x\ Re\ x < 1\ is-real\ x\ of-complex\ x \neq u$ 
  let  $?Mx = ?M\ (of-complex\ x)$ 
  have  $of-complex\ x \in unit-disc$ 
    using  $x\ cmod-eq-Re$ 
    by auto
  hence  $?Mx \in unit-disc$ 
    using  $\langle is-real\ a \rangle\ cmod\ a < 1 \rangle\ blaschke-unit-disc-fix[of\ a]$ 
    using unit-disc-fix-discI
    by blast
  hence  $?Mx \neq \infty_h$ 
    by auto
  moreover
  have  $of-complex\ x \in circline-set\ x-axis$ 
    using  $x$ 
    by auto
  hence  $?Mx \in circline-set\ x-axis$ 
    using blaschke-real-preserve-x-axis[OF \langle is-real\ a \rangle \langle cmod\ a < 1 \rangle, of\ of-complex\ x]
    by auto
  hence  $-1 < Re\ (to-complex\ ?Mx) \wedge Re\ (to-complex\ ?Mx) < 1 \wedge is-real\ (to-complex\ ?Mx)$ 
    using  $\langle ?Mx \neq \infty_h \rangle\ \langle ?Mx \in unit-disc \rangle$ 
    unfolding circline-set-x-axis
    by (auto simp add: cmod-eq-Re)
  moreover
  have  $?Mx \neq ?Mu$ 
    using  $\langle of-complex\ x \neq u \rangle$ 
    by simp

```

```

ultimately
have ?P ?Mx ?Mu
  using *[rule-format, of to-complex ?Mx] ⟨?Mx ≠ ∞h⟩
  by simp
then obtain Mz where
  ?R Mz ?Mx ?Mu
  by blast
have of-complex Mz ∈ unit-disc of-complex Mz ∈ circline-set x-axis
  using ⟨?R Mz ?Mx ?Mu⟩
  using cmod-eq-Re
  by auto

let ?Minv = - (blaschke a)
let ?z = moebius-pt ?Minv (of-complex Mz)
have ?z ∈ unit-disc
  using ⟨of-complex Mz ∈ unit-disc⟩ ⟨cmod a < 1⟩
  by auto
moreover
have ?z ∈ circline-set x-axis
  using ⟨of-complex Mz ∈ circline-set x-axis⟩
  using blaschke-real-preserve-x-axis ⟨is-real a⟩ ⟨cmod a < 1⟩
  by fastforce
ultimately
have z1: -1 < Re (to-complex ?z) Re (to-complex ?z) < 1 is-real (to-complex ?z)
  using inf-or-of-complex[of ?z]
  unfolding circline-set-x-axis
  by (auto simp add: cmod-eq-Re)

have z2: poincare-distance ?z (of-complex x) = poincare-distance ?z u
  using ⟨?R Mz ?Mx ?Mu⟩ ⟨cmod a < 1⟩ ⟨?z ∈ unit-disc⟩ ⟨of-complex x ∈ unit-disc⟩ ⟨u ∈ unit-disc⟩
by (metis blaschke-preserve-distance-formula blaschke-unit-disc-fix moebius-pt-comp-inv-right poincare-distance-formula
uminus-moebius-def unit-disc-fix-discI unit-disc-iff-cmod-lt-1)
show ?P (of-complex x) u
proof
  show ?R (to-complex ?z) (of-complex x) u
    using z1 z2 ⟨?z ∈ unit-disc⟩ inf-or-of-complex[of ?z]
    by auto
next
fix z'
assume ?R z' (of-complex x) u
hence of-complex z' ∈ unit-disc of-complex z' ∈ circline-set x-axis
  by (auto simp add: cmod-eq-Re)
let ?Mz' = ?M (of-complex z')
have ?Mz' ∈ unit-disc ?Mz' ∈ circline-set x-axis
  using ⟨of-complex z' ∈ unit-disc⟩ ⟨of-complex z' ∈ circline-set x-axis⟩ ⟨cmod a < 1⟩ ⟨is-real a⟩
  using blaschke-unit-disc-fix unit-disc-fix-discI
  using blaschke-real-preserve-x-axis circline-set-x-axis
  by blast+
hence -1 < Re (to-complex ?Mz') Re (to-complex ?Mz') < 1 is-real (to-complex ?Mz')
  unfolding circline-set-x-axis
  by (auto simp add: cmod-eq-Re)
moreover
have poincare-distance ?Mz' ?Mx = poincare-distance ?Mz' ?Mu
  using ⟨?R z' (of-complex x) u⟩
  using ⟨cmod a < 1⟩ ⟨of-complex x ∈ unit-disc⟩ ⟨of-complex z' ∈ unit-disc⟩ ⟨u ∈ unit-disc⟩
  by auto
ultimately
have ?R (to-complex ?Mz') ?Mx ?Mu
  using ⟨?Mz' ∈ unit-disc⟩ inf-or-of-complex[of ?Mz']
  by auto
hence ?Mz' = of-complex Mz
  using ⟨?P ?Mx ?Mu⟩ ⟨?R Mz ?Mx ?Mu⟩
by (metis ⟨moebius-pt (blaschke a) (of-complex z') ∈ unit-disc⟩ ⟨of-complex Mz ∈ unit-disc⟩ to-complex-of-complex
unit-disc-to-complex-inj)
thus z' = to-complex ?z
  by (simp add: moebius-pt-invert)

```



```

qed
qed
qed
thus ?thesis
  using assms
  by (metis to-complex-of-complex)
qed

```

5.3 Triangle inequality

lemma *poincare-distance-formula-zero-sum*:

assumes $u \in \text{unit-disc}$ and $v \in \text{unit-disc}$

shows $\text{poincare-distance } u \ 0_h + \text{poincare-distance } 0_h \ v =$

(let $u' = \text{cmod } (\text{to-complex } u)$; $v' = \text{cmod } (\text{to-complex } v)$
in $\text{arcosh } (((1 + u'^2) * (1 + v'^2) + 4 * u' * v') / ((1 - u'^2) * (1 - v'^2))))$)

proof–

obtain $u' \ v'$ **where** $uw: u' = \text{to-complex } u \ v' = \text{to-complex } v$

by *auto*

have $uv': u = \text{of-complex } u' \ v = \text{of-complex } v'$

using $uw \ \text{assms } \text{inf-or-of-complex}[of \ u] \ \text{inf-or-of-complex}[of \ v]$

by *auto*

let $?u' = \text{cmod } u' \ \text{and} \ ?v' = \text{cmod } v'$

have $\text{disc}: ?u'^2 < 1 \ ?v'^2 < 1$

using $\text{unit-disc-cmod-square-lt-1}[OF \ \langle u \in \text{unit-disc} \rangle]$

using $\text{unit-disc-cmod-square-lt-1}[OF \ \langle v \in \text{unit-disc} \rangle] \ uv$

by *auto*

thm *arcosh-add*

have $\text{arcosh } (1 + 2 * ?u'^2 / (1 - ?u'^2)) + \text{arcosh } (1 + 2 * ?v'^2 / (1 - ?v'^2)) =$

$\text{arcosh } (((1 + ?u'^2) * (1 + ?v'^2) + 4 * ?u' * ?v') / ((1 - ?u'^2) * (1 - ?v'^2)))$ (**is** $\text{arcosh } ?ll + \text{arcosh } ?rr =$
 $\text{arcosh } ?r$)

proof (*subst arcosh-add*)

show $?ll \geq 1 \ ?rr \geq 1$

using *disc*

by *auto*

next

show $\text{arcosh } (((1 + 2 * ?u'^2 / (1 - ?u'^2)) * (1 + 2 * ?v'^2 / (1 - ?v'^2)) +$

$\text{sqrt } (((1 + 2 * ?u'^2 / (1 - ?u'^2))^2 - 1) * ((1 + 2 * ?v'^2 / (1 - ?v'^2))^2 - 1))) =$

$\text{arcosh } ?r$ (**is** $\text{arcosh } ?l = -$)

proof–

have $1 + 2 * ?u'^2 / (1 - ?u'^2) = (1 + ?u'^2) / (1 - ?u'^2)$

using *disc*

by (*subst add-divide-eq-iff, simp-all*)

moreover

have $1 + 2 * ?v'^2 / (1 - ?v'^2) = (1 + ?v'^2) / (1 - ?v'^2)$

using *disc*

by (*subst add-divide-eq-iff, simp-all*)

moreover

have $\text{sqrt } (((1 + 2 * ?u'^2 / (1 - ?u'^2))^2 - 1) * ((1 + 2 * ?v'^2 / (1 - ?v'^2))^2 - 1)) =$

$(4 * ?u' * ?v') / ((1 - ?u'^2) * (1 - ?v'^2))$ (**is** $\text{sqrt } ?s = ?t$)

proof–

have $?s = ?t^2$

using *disc*

apply (*subst add-divide-eq-iff, simp*)+

apply (*subst power-divide*)+

apply *simp*

apply (*subst divide-diff-eq-iff, simp*)+

apply (*simp add: power2-eq-square field-simps*)

done

thus *?thesis*

using *disc*

by *simp*

qed

ultimately

have $?l = ?r$

```

    using disc
  by simp (subst add-divide-distrib, simp)
thus ?thesis
  by simp
qed
qed
thus ?thesis
  using w' assms
  using poincare-distance-formula
  by (simp add: Let-def)
qed

```

lemma poincare-distance-triangle-inequality:

```

assumes u ∈ unit-disc and v ∈ unit-disc and w ∈ unit-disc
shows poincare-distance u v + poincare-distance v w ≥ poincare-distance u w (is ?P' u v w)

```

proof—

```

have ∀ w. w ∈ unit-disc → ?P' u v w (is ?P v u)

```

```

proof (rule wlog-x-axis[where P=?P])

```

```

  fix x

```

```

  assume is-real x 0 ≤ Re x Re x < 1

```

```

  hence of-complex x ∈ unit-disc

```

```

    by (simp add: cmod-eq-Re)

```

```

show ?P 0h (of-complex x)

```

```

proof safe

```

```

  fix w

```

```

  assume w ∈ unit-disc

```

```

  then obtain w' where w = of-complex w'

```

```

    using inf-or-of-complex[of w]

```

```

    by auto

```

```

let ?x = cmod x and ?w = cmod w' and ?xw = cmod (x - w')

```

```

have disc: ?x2 < 1 ?w2 < 1

```

```

  using unit-disc-cmod-square-lt-1[OF ⟨of-complex x ∈ unit-disc⟩]

```

```

  using unit-disc-cmod-square-lt-1[OF ⟨w ∈ unit-disc⟩] w

```

```

  by auto

```

```

have poincare-distance (of-complex x) 0h + poincare-distance 0h w =

```

```

  arcosh (((1 + ?x2) * (1 + ?w2) + 4 * ?x * ?w) / ((1 - ?x2) * (1 - ?w2))) (is - = arcosh ?r1)

```

```

  using poincare-distance-formula-zero-sum[OF ⟨of-complex x ∈ unit-disc⟩ ⟨w ∈ unit-disc⟩] w

```

```

  by (simp add: Let-def)

```

moreover

```

have poincare-distance (of-complex x) (of-complex w') =

```

```

  arcosh (((1 - ?x2) * (1 - ?w2) + 2 * ?xw2) / ((1 - ?x2) * (1 - ?w2))) (is - = arcosh ?r2)

```

```

  using disc

```

```

  using poincare-distance-formula[OF ⟨of-complex x ∈ unit-disc⟩ ⟨w ∈ unit-disc⟩] w

```

```

  by (subst add-divide-distrib) simp

```

moreover

```

have *: (1 - ?x2) * (1 - ?w2) + 2 * ?xw2 ≤ (1 + ?x2) * (1 + ?w2) + 4 * ?x * ?w

```

proof—

```

  have (cmod (x - w'))2 ≤ (cmod x + cmod w')2

```

```

    using norm-triangle-ineq4[of x w']

```

```

    by (simp add: power-mono)

```

```

  thus ?thesis

```

```

    by (simp add: field-simps power2-sum)

```

qed

```

have arcosh ?r1 ≥ arcosh ?r2

```

```

proof (subst arcosh-mono)

```

```

  show ?r1 ≥ 1

```

```

    using disc

```

```

    by (smt * le-divide-eq-1-pos mult-pos-pos zero-le-power2)

```

next

```

  show ?r2 ≥ 1

```

```

    using disc

```

```

    by simp

```

```

next
  show ?r1 ≥ ?r2
  using disc
  using *
  by (subst divide-right-mono, simp-all)
qed
ultimately
show poincare-distance (of-complex x) w ≤ poincare-distance (of-complex x) 0_h + poincare-distance 0_h w
  using ⟨of-complex x ∈ unit-disc⟩ ⟨w ∈ unit-disc⟩ w
  using poincare-distance-formula
  by simp
qed
next
show v ∈ unit-disc u ∈ unit-disc
  by fact+
next
fix M u v
assume *: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc
assume **: ?P (moebius-pt M u) (moebius-pt M v)
show ?P u v
proof safe
  fix w
  assume w ∈ unit-disc
  thus ?P' v u w
    by (metis * ** unit-disc-fix-discI unit-disc-fix-preserve-poincare-distance)
qed
qed
thus ?thesis
  using assms
  by auto
qed

end
theory Poincare-Circles
  imports Poincare-Distance
begin

```

6 H-circles in the Poincaré model

Circles consist of points that are at the same distance from the center.

definition *poincare-circle* :: *complex-homo* ⇒ *real* ⇒ *complex-homo set* **where**
poincare-circle z r = {z'. z' ∈ unit-disc ∧ poincare-distance z z' = r}

Each h-circle in the Poincaré model is represented by an Euclidean circle in the model — the center and radius of that euclidean circle are determined by the following formulas.

definition *poincare-circle-euclidean* :: *complex-homo* ⇒ *real* ⇒ *euclidean-circle* **where**
poincare-circle-euclidean z r =
 (let R = (cosh r - 1) / 2;
 z' = to-complex z;
 cz = 1 - (cmod z')²;
 k = cz * R + 1
 in (z' / k, cz * sqrt(R * (R + 1)) / k))

That Euclidean circle has a positive radius and is always fully within the disc.

lemma *poincare-circle-in-disc*:

assumes r > 0 **and** z ∈ unit-disc **and** (ze, re) = *poincare-circle-euclidean* z r
shows cmod ze < 1 re > 0 ∀ x ∈ circle ze re. cmod x < 1

proof—

```

let ?R = (cosh r - 1) / 2
let ?z' = to-complex z
let ?cz = 1 - (cmod ?z')2
let ?k = ?cz * ?R + 1
let ?ze = ?z' / ?k
let ?re = ?cz * sqrt(?R * (?R + 1)) / ?k

```

```

from ⟨z ∈ unit-disc⟩
obtain z' where z': z = of-complex z'
  using inf-or-of-complex[of z]
  by auto

hence z' = ?z'
  by simp

obtain cz where cz: cz = (1 - (cmod z')2)
  by auto

have cz > 0 cz ≤ 1
  using ⟨z ∈ unit-disc⟩ z' cz
  using unit-disc-cmod-square-lt-1
  by fastforce+

obtain R where R: R = ?R
  by blast

have R > 0
  using cosh-gt-1[of r] ⟨r > 0⟩
  by (subst R) simp

obtain k where k: k = cz * R + 1
  by auto

have k > 1
  using k ⟨R > 0⟩ ⟨cz > 0⟩
  by simp

hence cmod k = k
  by simp

let ?RR = cz * sqrt(R * (R + 1)) / k

have cmod z' + cz * sqrt(R * (R + 1)) < k
proof-
  have ((R+1)-R)2 > 0
    by simp
  hence (R+1)2 - 2*R*(R+1) + R2 > 0
    unfolding power2-diff
    by (simp add: field-simps)
  hence (R+1)2 + 2*R*(R+1) + R2 - 4*R*(R+1) > 0
    by simp
  hence (2*R+1)2 / 4 > R*(R+1)
    using power2-sum[of R+1 R]
    by (simp add: field-simps)
  hence sqrt(R*(R+1)) < (2*R+1) / 2
    using ⟨R > 0⟩
    by (smt arith-geo-mean-sqrt power-divide real-sqrt-four real-sqrt-pow2 zero-le-mult-iff)
  hence sqrt(R*(R+1)) - R < 1/2
    by (simp add: field-simps)
  hence (1 + (cmod z')) * (sqrt(R*(R+1)) - R) < (1 + (cmod z')) * (1 / 2)
    by (subst mult-strict-left-mono, simp, smt norm-not-less-zero, simp)
  also have ... < 1
    using ⟨z ∈ unit-disc⟩ z'
    by auto
  finally have (1 - cmod z') * ((1 + cmod z') * (sqrt(R*(R+1)) - R)) < (1 - cmod z') * 1
    using ⟨z ∈ unit-disc⟩ z'
    by (subst mult-strict-left-mono, simp-all)
  hence cz * (sqrt(R*(R+1)) - R) < 1 - cmod z'
    using square-diff-square-factored[of 1 cmod z']
    by (subst cz, subst (asm) mult.assoc[symmetric], simp add: power2-eq-square field-simps)
  hence cmod z' + cz * sqrt(R*(R+1)) < 1 + R * cz
    by (simp add: field-simps)

```

```

thus ?thesis
  using k
  by (simp add: field-simps)
qed
hence cmod z' / k + cz * sqrt(R * (R + 1)) / k < 1
  using ⟨k > 1⟩
  unfolding add-divide-distrib[symmetric]
  by simp
hence cmod (z' / k) + cz * sqrt(R * (R + 1)) / k < 1
  using ⟨k > 1⟩
  by simp
hence cmod ?ze + ?re < 1
  using k cz ⟨R = ?R⟩ z'
  by simp

```

moreover

```

have cz * sqrt(R * (R + 1)) / k > 0
  using ⟨cz > 0⟩ ⟨R > 0⟩ ⟨k > 1⟩
  by auto
hence ?re > 0
  using k cz ⟨R = ?R⟩ z'
  by simp

```

moreover

```

have cmod ?ze < 1
  using ⟨cmod ?ze + ?re < 1⟩ ⟨?re > 0⟩
  by simp

```

moreover

```

have ze = ?ze re = ?re
  using ⟨(ze, re) = poincare-circle-euclidean z r⟩
  unfolding poincare-circle-euclidean-def Let-def
  by simp-all

```

moreover

```

have ∀ x ∈ circle ze re. cmod x ≤ cmod ze + re
  using norm-triangle-ineq2[of - ze]
  unfolding circle-def
  by (smt mem-Collect-eq)

```

ultimately

```

show cmod ze < 1 re > 0 ∀ x ∈ circle ze re. cmod x < 1
  by auto

```

qed

The connection between the points on the h-circle and its corresponding Euclidean circle.

lemma poincare-circle-is-euclidean-circle:

assumes z ∈ unit-disc **and** r > 0

shows let (Ze, Re) = poincare-circle-euclidean z r
in of-complex ' (circle Ze Re) = poincare-circle z r

proof –

```

{
  fix x
  let ?z = to-complex z

  from assms obtain z' where z' : z = of-complex z' cmod z' < 1
    using inf-or-of-complex[of z]
    by auto

```

```

have *: ∧ x. cmod x < 1 ⇒ 1 - (cmod x)2 > 0
by (metis less-iff-diff-less-0 minus-diff-eq mult.left-neutral neg-less-0-iff-less norm-mult-less norm-power power2-eq-square)

```

```

let ?R = (cosh r - 1) / 2
obtain R where R: R = ?R
  by blast

let ?cx = 1 - (cmod x)2 and ?cz = 1 - (cmod z')2 and ?czx = (cmod (z' - x))2

let ?k = 1 + R * ?cz
obtain k where k: k = ?k
  by blast
have R > 0
  using R cosh-gt-1[OF ‹r > 0›]
  by simp

hence k > 1
  using assms z' k *[of z']
  by auto
hence **: cor k ≠ 0
  by (smt of-real-eq-0-iff)

have of-complex x ∈ poincare-circle z r ↔ cmod x < 1 ∧ poincare-distance z (of-complex x) = r
  unfolding poincare-circle-def
  by auto
also have ... ↔ cmod x < 1 ∧ poincare-distance-formula' ?z x = cosh r
  using poincare-distance-formula[of z of-complex x] cosh-dist[of z of-complex x]
  unfolding poincare-distance-formula-def
  using assms
  using arcosh-cosh-real
  by auto
also have ... ↔ cmod x < 1 ∧ ?czx / (?cz * ?cx) = ?R
  using z'
  by (simp add: field-simps)
also have ... ↔ cmod x < 1 ∧ ?czx = ?R * ?cx * ?cz
  using assms z' *[of z'] *[of x]
  using nonzero-divide-eq-eq[of (1 - (cmod x)2) * (1 - (cmod z')2) (cmod (z' - x))2 ?R]
  by (auto, simp add: field-simps)
also have ... ↔ cmod x < 1 ∧ (z' - x) * (cnj z' - cnj x) = R * ?cz * (1 - x * cnj x) (is - ↔ - ∧ ?l = ?r)
proof-
  let ?l = (z' - x) * (cnj z' - cnj x) and ?r = R * (1 - Re (z' * cnj z')) * (1 - x * cnj x)
  have is-real ?l
    using eq-cnj-iff-real[of ?l]
    by simp
  moreover
  have is-real ?r
    using eq-cnj-iff-real[of 1 - x * cnj x]
    using Im-complex-of-real[of R * (1 - Re (z' * cnj z'))]
    by simp
  ultimately
  show ?thesis
    apply (subst R[symmetric])
    apply (subst cmod-square)+
    apply (subst complex-eq-iff-Re-eq, simp-all add: field-simps)
    done
qed
also have ... ↔ cmod x < 1 ∧ z' * cnj z' - x * cnj z' - cnj x * z' + x * cnj x = R * ?cz - R * ?cz * x * cnj x
  unfolding right-diff-distrib left-diff-distrib
  by (simp add: field-simps)
also have ... ↔ cmod x < 1 ∧ k * (x * cnj x) - x * cnj z' - cnj x * z' + z' * cnj z' = R * ?cz (is - ↔ - ∧
?lhs = ?rhs)
  by (subst k) (auto simp add: field-simps)
also have ... ↔ cmod x < 1 ∧ (k * x * cnj x - x * cnj z' - cnj x * z' + z' * cnj z') / k = (R * ?cz) / k
  using **
  by (auto simp add: Groups.mult-ac(1))
also have ... ↔ cmod x < 1 ∧ x * cnj x - x * cnj z' / k - cnj x * z' / k + z' * cnj z' / k = (R * ?cz) / k
  using **

```

```

    unfolding add-divide-distrib diff-divide-distrib
  by auto
also have ...  $\longleftrightarrow$   $cmod\ x < 1 \wedge (x - z'/k) * cnj(x - z'/k) = (R * ?cz) / k + (z' / k) * cnj(z' / k) - z' * cnj\ z' / k$ 
  by (auto simp add: field-simps diff-divide-distrib)
also have ...  $\longleftrightarrow$   $cmod\ x < 1 \wedge (cmod\ (x - z'/k))^2 = (R * ?cz) / k + (cmod\ z')^2 / k^2 - (cmod\ z')^2 / k$ 
  apply (subst complex-mult-cnj-cmod)+
  apply (subst complex-eq-if-Re-eq)
  apply (simp-all add: power-divide)
done
also have ...  $\longleftrightarrow$   $cmod\ x < 1 \wedge (cmod\ (x - z'/k))^2 = (R * ?cz * k + (cmod\ z')^2 - (cmod\ z')^2 * k) / k^2$ 
  using **
  unfolding add-divide-distrib diff-divide-distrib
  by (simp add: power2-eq-square)
also have ...  $\longleftrightarrow$   $cmod\ x < 1 \wedge (cmod\ (x - z'/k))^2 = ?cz^2 * R * (R + 1) / k^2$  (is -  $\longleftrightarrow$  -  $\wedge$   $?a^2 = ?b$ )
proof-
  have *:  $R * (1 - (cmod\ z')^2) * k + (cmod\ z')^2 - (cmod\ z')^2 * k = (1 - (cmod\ z')^2)^2 * R * (R + 1)$ 
    by (subst k)+ (simp add: field-simps power2-diff)
  thus ?thesis
    by (subst *, simp)
qed
also have ...  $\longleftrightarrow$   $cmod\ x < 1 \wedge cmod\ (x - z'/k) = ?cz * sqrt\ (R * (R + 1)) / k$ 
  using  $\langle R > 0 \rangle * [of\ z'] ** \langle k > 1 \rangle \langle z \in unit-disc \rangle z'$ 
  using real-sqrt-unique[ $of\ ?a\ ?b$ , symmetric]
  by (auto simp add: real-sqrt-divide real-sqrt-mult power-divide power-mult-distrib)
finally
have of-complex  $x \in poincare-circle\ z\ r \longleftrightarrow cmod\ x < 1 \wedge x \in circle\ (z'/k)\ (?cz * sqrt\ (R * (R+1)) / k)$ 
  unfolding circle-def  $z'\ k\ R$ 
  by simp
hence of-complex  $x \in poincare-circle\ z\ r \longleftrightarrow (let\ (Ze, Re) = poincare-circle-euclidean\ z\ r\ in\ cmod\ x < 1 \wedge x \in circle\ Ze\ Re)$ 
  unfolding poincare-circle-euclidean-def Let-def circle-def
  using  $z'\ R\ k$ 
  by (simp add: field-simps)
hence of-complex  $x \in poincare-circle\ z\ r \longleftrightarrow (let\ (Ze, Re) = poincare-circle-euclidean\ z\ r\ in\ x \in circle\ Ze\ Re)$ 
  using poincare-circle-in-disc[ $OF\ \langle r > 0 \rangle \langle z \in unit-disc \rangle$ ]
  by auto
} note * = this
show ?thesis
  unfolding Let-def
proof safe
  fix  $Ze\ Re\ x$ 
  assume poincare-circle-euclidean  $z\ r = (Ze, Re)\ x \in circle\ Ze\ Re$ 
  thus of-complex  $x \in poincare-circle\ z\ r$ 
    using *[ $of\ x$ ]
    by simp
next
  fix  $Ze\ Re\ x$ 
  assume **: poincare-circle-euclidean  $z\ r = (Ze, Re)\ x \in poincare-circle\ z\ r$ 
  then obtain  $x'$  where  $x' : x = of-complex\ x'$ 
    unfolding poincare-circle-def
    using inf-or-of-complex[ $of\ x$ ]
    by auto
  hence  $x' \in circle\ Ze\ Re$ 
    using *[ $of\ x'$ ] **
    by simp
  thus  $x \in of-complex\ \langle circle\ Ze\ Re$ 
    using  $x'$ 
    by auto
qed
qed

```

6.1 Intersection of circles in special positions

Two h-circles centered at the x-axis intersect at mutually conjugate points

lemma intersect-poincare-circles-x-axis:

assumes z : *is-real* $z1$ **and** *is-real* $z2$ **and** $r1 > 0$ **and** $r2 > 0$ **and**
 $-1 < \text{Re } z1$ **and** $\text{Re } z1 < 1$ **and** $-1 < \text{Re } z2$ **and** $\text{Re } z2 < 1$ **and**
 $z1 \neq z2$
assumes $x1$: $x1 \in \text{poincare-circle (of-complex } z1) r1 \cap \text{poincare-circle (of-complex } z2) r2$ **and**
 $x2$: $x2 \in \text{poincare-circle (of-complex } z1) r1 \cap \text{poincare-circle (of-complex } z2) r2$ **and**
 $x1 \neq x2$

shows $x1 = \text{conjugate } x2$

proof–

have *in-disc*: $\text{of-complex } z1 \in \text{unit-disc of-complex } z2 \in \text{unit-disc}$
using *assms*
by (*auto simp add: cmod-eq-Re*)

obtain $x1' x2'$ **where** x' : $x1 = \text{of-complex } x1' x2 = \text{of-complex } x2'$
using $x1 x2$
using *inf-or-of-complex[of x1] inf-or-of-complex[of x2]*
unfolding *poincare-circle-def*
by *auto*

obtain $Ze1 Re1$ **where** 1 : $(Ze1, Re1) = \text{poincare-circle-euclidean (of-complex } z1) r1$
by (*metis poincare-circle-euclidean-def*)

obtain $Ze2 Re2$ **where** 2 : $(Ze2, Re2) = \text{poincare-circle-euclidean (of-complex } z2) r2$
by (*metis poincare-circle-euclidean-def*)

have *circle*: $x1' \in \text{circle } Ze1 Re1 \cap \text{circle } Ze2 Re2$ $x2' \in \text{circle } Ze1 Re1 \cap \text{circle } Ze2 Re2$
using *poincare-circle-is-euclidean-circle[of of-complex z1 r1]*
using *poincare-circle-is-euclidean-circle[of of-complex z2 r2]*
using *assms 1 2 <of-complex z1 ∈ unit-disc> <of-complex z2 ∈ unit-disc> x'*
by *auto (metis image-iff of-complex-inj)+*

have *is-real* $Ze1$ *is-real* $Ze2$
using $1 2$ *<is-real z1> <is-real z2>*
by (*simp-all add: poincare-circle-euclidean-def Let-def*)

have $Re1 > 0$ $Re2 > 0$
using $1 2$ *in-disc <r1 > 0> <r2 > 0>*
using *poincare-circle-in-disc(2)[of r1 of-complex z1 Ze1 Re1]*
using *poincare-circle-in-disc(2)[of r2 of-complex z2 Ze2 Re2]*
by *auto*

have $Ze1 \neq Ze2$

proof (*rule ccontr*)
assume $\neg ?thesis$
hence *eq*: $Ze1 = Ze2$ $Re1 = Re2$
using *circle(1)*
unfolding *circle-def*
by *auto*

let $?A = Ze1 - Re1$ **and** $?B = Ze1 + Re1$

have $?A \in \text{circle } Ze1 Re1$ $?B \in \text{circle } Ze1 Re1$

using *<Re1 > 0>*
unfolding *circle-def*
by *simp-all*

hence *of-complex ?A ∈ poincare-circle (of-complex z1) r1 of-complex ?B ∈ poincare-circle (of-complex z1) r1*
of-complex ?A ∈ poincare-circle (of-complex z2) r2 of-complex ?B ∈ poincare-circle (of-complex z2) r2

using *eq*
using *poincare-circle-is-euclidean-circle[OF <of-complex z1 ∈ unit-disc> <r1 > 0>]*
using *poincare-circle-is-euclidean-circle[OF <of-complex z2 ∈ unit-disc> <r2 > 0>]*
using $1 2$
by *auto blast+*

hence *poincare-distance (of-complex z1) (of-complex ?A) = poincare-distance (of-complex z1) (of-complex ?B)*
poincare-distance (of-complex z2) (of-complex ?A) = poincare-distance (of-complex z2) (of-complex ?B)
 $-1 < \text{Re } (Ze1 - Re1)$ $\text{Re } (Ze1 - Re1) < 1$ $-1 < \text{Re } (Ze1 + Re1)$ $\text{Re } (Ze1 + Re1) < 1$

using *<is-real Ze1> <is-real Ze2>*
unfolding *poincare-circle-def*
by (*auto simp add: cmod-eq-Re*)

hence $z1 = z2$
using *unique-midpoint-x-axis[of Ze1 - Re1 Ze1 + Re1]*


```

    using ‹is-real Ze1› ‹is-real z1› ‹is-real z2› ‹Re1 > 0› ‹-1 < Re z1› ‹Re z1 < 1› ‹-1 < Re z2› ‹Re z2 < 1›
  by auto
  thus False
    using ‹z1 ≠ z2›
  by simp
qed

```

```

hence *: (Re x1')^2 + (Im x1')^2 - 2 * Re x1' * Ze1 + Ze1 * Ze1 - cor (Re1 * Re1) = 0
        (Re x1')^2 + (Im x1')^2 - 2 * Re x1' * Ze2 + Ze2 * Ze2 - cor (Re2 * Re2) = 0
        (Re x2')^2 + (Im x2')^2 - 2 * Re x2' * Ze1 + Ze1 * Ze1 - cor (Re1 * Re1) = 0
        (Re x2')^2 + (Im x2')^2 - 2 * Re x2' * Ze2 + Ze2 * Ze2 - cor (Re2 * Re2) = 0
  using circle-equation[of Re1 Ze1] circle-equation[of Re2 Ze2] circle
  using eq-cn timer-iff-real[of Ze1] ‹is-real Ze1› ‹Re1 > 0›
  using eq-cn timer-iff-real[of Ze2] ‹is-real Ze2› ‹Re2 > 0›
  using complex-add-cn timer[of x1'] complex-add-cn timer[of x2']
  using distrib-left[of Ze1 x1' cnj x1'] distrib-left[of Ze2 x1' cnj x1']
  using distrib-left[of Ze1 x2' cnj x2'] distrib-left[of Ze2 x2' cnj x2']
  by (auto simp add: complex-mult-cn timer power2-eq-square field-simps)

```

```

hence - 2 * Re x1' * Ze1 + Ze1 * Ze1 - cor (Re1 * Re1) = - 2 * Re x1' * Ze2 + Ze2 * Ze2 - cor (Re2 * Re2)
      - 2 * Re x2' * Ze1 + Ze1 * Ze1 - cor (Re1 * Re1) = - 2 * Re x2' * Ze2 + Ze2 * Ze2 - cor (Re2 * Re2)
  by (smt add-diff-cancel-right' add-diff-eq eq-iff-diff-eq-0 minus-diff-eq mult-minus-left of-real-minus)+
hence 2 * Re x1' * (Ze2 - Ze1) = (Ze2 * Ze2 - cor (Re2 * Re2)) - (Ze1 * Ze1 - cor (Re1 * Re1))
      2 * Re x2' * (Ze2 - Ze1) = (Ze2 * Ze2 - cor (Re2 * Re2)) - (Ze1 * Ze1 - cor (Re1 * Re1))
  by simp-all (simp add: field-simps)+
hence 2 * Re x1' * (Ze2 - Ze1) = 2 * Re x2' * (Ze2 - Ze1)
  by simp
hence Re x1' = Re x2'
  using ‹Ze1 ≠ Ze2›
  by simp
moreover
hence (Im x1')^2 = (Im x2')^2
  using *(1) *(3)
  by (simp add: ‹is-real Ze1› complex-eq-if-Re-eq)
hence Im x1' = Im x2' ∨ Im x1' = -Im x2'
  using power2-eq-iff
  by blast
ultimately
show ?thesis
  using x' ‹x1 ≠ x2›
  using complex.expand
  by (metis cnj code complex-surj conjugate-of-complex)
qed

```

Two h-circles of the same radius centered at mutually conjugate points intersect at the x-axis

```

lemma intersect-poincare-circles-conjugate-centers:
  assumes in-disc: z1 ∈ unit-disc z2 ∈ unit-disc and
          z1 ≠ z2 and z1 = conjugate z2 and r > 0 and
          u: u ∈ poincare-circle z1 r ∩ poincare-circle z2 r
  shows is-real (to-complex u)

```

```

proof -
  obtain z1e r1e z2e r2e where
    euclidean: (z1e, r1e) = poincare-circle-euclidean z1 r
              (z2e, r2e) = poincare-circle-euclidean z2 r
  by (metis poincare-circle-euclidean-def)
  obtain z1' z2' where z': z1 = of-complex z1' z2 = of-complex z2'
  using inf-or-of-complex[of z1] inf-or-of-complex[of z2] in-disc
  by auto
  obtain u' where u': u = of-complex u'
  using u inf-or-of-complex[of u]
  by (auto simp add: poincare-circle-def)
  have z1' = cnj z2'
  using ‹z1 = conjugate z2› z'
  by (auto simp add: of-complex-inj)
  moreover
  let ?cz = 1 - (cmod z2')^2

```

```

let ?den = ?cz * (cosh r - 1) / 2 + 1
have ?cz > 0
  using in-disc z'
  by (simp add: cmod-def)
hence ?den ≥ 1
  using cosh-gt-1[OF ‹r > 0›]
  by auto
hence ?den ≠ 0
  by simp
hence cor ?den ≠ 0
  using of-real-eq-0-iff
  by blast
ultimately
have r1e = r2e z1e = cnj z2e z1e ≠ z2e
  using z' euclidean ‹z1 ≠ z2›
  unfolding poincare-circle-euclidean-def Let-def
  by simp-all metis

hence u' ∈ circle (cnj z2e) r2e ∩ circle z2e r2e z2e ≠ cnj z2e
  using euclidean u u'
  using poincare-circle-is-euclidean-circle[of z1 r]
  using poincare-circle-is-euclidean-circle[of z2 r]
  using in-disc ‹r > 0›
  by auto (metis image-iff of-complex-inj)+
hence (cmod (u' - z2e))2 = (cmod(u' - cnj z2e))2
  by (simp add: circle-def)
hence (u' - z2e) * (cnj u' - cnj z2e) = (u' - cnj z2e) * (cnj u' - z2e)
  by (metis complex-cnj-cnj complex-cnj-diff complex-norm-square)
hence (z2e - cnj z2e) * (u' - cnj u') = 0
  by (simp add: field-simps)
thus ?thesis
  using u' ‹z2e ≠ cnj z2e› eq-cnj-iff-real[of u']
  by simp
qed

```

6.2 Congruent triangles

For every pair of triangles such that its three pairs of sides are pairwise equal there is an h-isometry (a unit disc preserving Möbius transform, eventually composed with a conjugation) that maps one triangle onto the other.

lemma *unit-disc-fix-f-congruent-triangles*:

assumes

in-disc: $u \in \text{unit-disc } v \in \text{unit-disc } w \in \text{unit-disc}$ **and**
in-disc': $u' \in \text{unit-disc } v' \in \text{unit-disc } w' \in \text{unit-disc}$ **and**
d: *poincare-distance* $u v = \text{poincare-distance } u' v'$
poincare-distance $v w = \text{poincare-distance } v' w'$
poincare-distance $u w = \text{poincare-distance } u' w'$

shows

$\exists M. \text{unit-disc-fix-f } M \wedge M u = u' \wedge M v = v' \wedge M w = w'$

proof (cases $u = v \vee u = w \vee v = w$)

case *True*

thus ?thesis

using *assms*

using *poincare-distance-eq-0-iff*[of $u' v'$]

using *poincare-distance-eq-0-iff*[of $v' w'$]

using *poincare-distance-eq-0-iff*[of $u' w'$]

using *poincare-distance-eq-ex-moebius*[of $v w v' w'$]

using *poincare-distance-eq-ex-moebius*[of $u w u' w'$]

using *poincare-distance-eq-ex-moebius*[of $u v u' v'$]

by (metis *unit-disc-fix-f-def*)

next

case *False*

have $\forall w u' v' w'. w \in \text{unit-disc} \wedge u' \in \text{unit-disc} \wedge v' \in \text{unit-disc} \wedge w' \in \text{unit-disc} \wedge w \neq u \wedge w \neq v \wedge$
poincare-distance $u v = \text{poincare-distance } u' v' \wedge$
poincare-distance $v w = \text{poincare-distance } v' w' \wedge$

$\text{poincare-distance } u \ w = \text{poincare-distance } u' \ w' \longrightarrow$
 $(\exists M. \text{unit-disc-fix-f } M \wedge M \ u = u' \wedge M \ v = v' \wedge M \ w = w') \text{ (is } ?P \ u \ v)$
proof (*rule wlog-positive-x-axis*[**where** $P = ?P$])
show $v \in \text{unit-disc } u \in \text{unit-disc}$
by *fact+*
next
show $u \neq v$
using *False*
by *simp*
next
fix x
assume $x: \text{is-real } x \ 0 < \text{Re } x \ \text{Re } x < 1$

hence *of-complex* $x \neq 0_h$
using *of-complex-zero-iff*[*of* x]
by (*auto simp add: complex.expand*)

show $?P \ 0_h \ (\text{of-complex } x)$
proof *safe*
fix $w \ u' \ v' \ w'$
assume *in-disc*: $w \in \text{unit-disc } u' \in \text{unit-disc } v' \in \text{unit-disc } w' \in \text{unit-disc}$
assume *poincare-distance* $0_h \ (\text{of-complex } x) = \text{poincare-distance } u' \ v'$
then obtain M' **where** $M': \text{unit-disc-fix } M' \ \text{moebius-pt } M' \ u' = 0_h \ \text{moebius-pt } M' \ v' = (\text{of-complex } x)$
using *poincare-distance-eq-ex-moebius*[*of* $u' \ v' \ 0_h \ \text{of-complex } x$] *in-disc* x
by (*auto simp add: cmod-eq-Re*)

let $?w = \text{moebius-pt } M' \ w'$
have $?w \in \text{unit-disc}$
using $\langle \text{unit-disc-fix } M' \rangle \ \langle w' \in \text{unit-disc} \rangle$
by *simp*

assume $w \neq 0_h \ w \neq \text{of-complex } x$
hence *dist-gt-0*: *poincare-distance* $0_h \ w > 0$ *poincare-distance* (*of-complex* x) $w > 0$
using *poincare-distance-eq-0-iff*[*of* $0_h \ w$] *in-disc* *poincare-distance-ge0*[*of* $0_h \ w$]
using *poincare-distance-eq-0-iff*[*of* *of-complex* $x \ w$] *in-disc* *poincare-distance-ge0*[*of* *of-complex* $x \ w$]
using x
by (*simp-all add: cmod-eq-Re*)

assume *poincare-distance* (*of-complex* x) $w = \text{poincare-distance } v' \ w'$
poincare-distance $0_h \ w = \text{poincare-distance } u' \ w'$
hence *poincare-distance* $0_h \ ?w = \text{poincare-distance } 0_h \ w$
poincare-distance (*of-complex* x) $?w = \text{poincare-distance} \ (\text{of-complex } x) \ w$
using $M'(1) \ M'(2)$ [*symmetric*] $M'(3)$ [*symmetric*] *in-disc*
using *unit-disc-fix-preserve-poincare-distance*[*of* $M' \ u' \ w'$]
using *unit-disc-fix-preserve-poincare-distance*[*of* $M' \ v' \ w'$]
by *simp-all*
hence $?w \in \text{poincare-circle } 0_h \ (\text{poincare-distance } 0_h \ w) \cap \text{poincare-circle} \ (\text{of-complex } x) \ (\text{poincare-distance} \ (\text{of-complex } x) \ w)$
 $w \in \text{poincare-circle } 0_h \ (\text{poincare-distance } 0_h \ w) \cap \text{poincare-circle} \ (\text{of-complex } x) \ (\text{poincare-distance} \ (\text{of-complex } x) \ w)$
using $\langle ?w \in \text{unit-disc} \rangle \ \langle w \in \text{unit-disc} \rangle$
unfolding *poincare-circle-def*
by *simp-all*
hence $?w = w \vee ?w = \text{conjugate } w$
using *intersect-poincare-circles-x-axis*[*of* $0 \ x \ \text{poincare-distance } 0_h \ w \ \text{poincare-distance} \ (\text{of-complex } x) \ w \ ?w \ w$] x
using $\langle \text{of-complex } x \neq 0_h \rangle \ \text{dist-gt-0}$
using *poincare-distance-eq-0-iff*
by *auto*
thus $\exists M. \text{unit-disc-fix-f } M \wedge M \ 0_h = u' \wedge M \ (\text{of-complex } x) = v' \wedge M \ w = w'$
proof
assume *moebius-pt* $M' \ w' = w$
thus *?thesis*
using M'
using *moebius-pt-invert*[*of* $M' \ u' \ 0_h$]
using *moebius-pt-invert*[*of* $M' \ v' \ \text{of-complex } x$]
using *moebius-pt-invert*[*of* $M' \ w' \ w$]

```

    apply (rule-tac x=moebius-pt (-M') in exI)
    apply (simp add: unit-disc-fix-f-def)
    apply (rule-tac x=-M' in exI, simp)
    done
  next
    let ?M = moebius-pt (-M') ∘ conjugate
    assume moebius-pt M' w' = conjugate w
    hence ?M w = w'
      using moebius-pt-invert[of M' w' conjugate w]
      by simp
    moreover
    have ?M 0h = u' ?M (of-complex x) = v'
      using moebius-pt-invert[of M' u' 0h]
      using moebius-pt-invert[of M' v' of-complex x]
      using M' ⟨is-real x⟩ eq-cnj-iff-real[of x]
      by simp-all
    moreover
    have unit-disc-fix-f ?M
      using ⟨unit-disc-fix M'⟩
      unfolding unit-disc-fix-f-def
      by (rule-tac x=-M' in exI, simp)
    ultimately
    show ?thesis
      by blast
  qed
qed
next
  fix M u v
  assume 1: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc
  let ?Mu = moebius-pt M u and ?Mv = moebius-pt M v
  assume 2: ?P ?Mu ?Mv
  show ?P u v
  proof safe
    fix w u' v' w'
    let ?Mw = moebius-pt M w and ?Mu' = moebius-pt M u' and ?Mv' = moebius-pt M v' and ?Mw' = moebius-pt
M w'
    assume w ∈ unit-disc u' ∈ unit-disc v' ∈ unit-disc w' ∈ unit-disc w ≠ u w ≠ v
      poincare-distance u v = poincare-distance u' v'
      poincare-distance v w = poincare-distance v' w'
      poincare-distance u w = poincare-distance u' w'
    then obtain M' where M': unit-disc-fix-f M' M' ?Mu = ?Mu' M' ?Mv = ?Mv' M' ?Mw = ?Mw'
      using 1 2[rule-format, of ?Mw ?Mu' ?Mv' ?Mw']
      by auto

    let ?M = moebius-pt (-M) ∘ M' ∘ moebius-pt M
    show ∃ M. unit-disc-fix-f M ∧ M u = u' ∧ M v = v' ∧ M w = w'
    proof (rule-tac x=?M in exI, safe)
      show unit-disc-fix-f ?M
        using M'(1) ⟨unit-disc-fix M⟩
        by (subst unit-disc-fix-f-comp, subst unit-disc-fix-f-comp, simp-all)
    next
      show ?M u = u' ?M v = v' ?M w = w'
        using M'
        by auto
    qed
  qed
qed
thus ?thesis
  using assms False
  by auto
qed
end
theory Poincare-Between
  imports Poincare-Distance
begin

```

7 H-betweenness in the Poincaré model

7.1 H-betweenness expressed by a cross-ratio

The point v is h-between u and w if the cross-ratio between the pairs u and w and v and inverse of v is real and negative.

definition *poincare-between* :: *complex-homo* \Rightarrow *complex-homo* \Rightarrow *complex-homo* \Rightarrow *bool* **where**
poincare-between $u\ v\ w \iff$
 $u = v \vee v = w \vee$
 $(\text{let } cr = \text{cross-ratio } u\ v\ w\ (\text{inversion } v)$
 $\text{in } \text{is-real } (\text{to-complex } cr) \wedge \text{Re } (\text{to-complex } cr) < 0)$

7.1.1 H-betweenness is preserved by h-isometries

Since they preserve cross-ratio and inversion, h-isometries (unit disc preserving Möbius transformations and conjugation) preserve h-betweenness.

lemma *unit-disc-fix-moebius-preserve-poincare-between* [*simp*]:
assumes *unit-disc-fix* M **and** $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$ **and** $w \in \text{unit-disc}$
shows *poincare-between* (*moebius-pt* $M\ u$) (*moebius-pt* $M\ v$) (*moebius-pt* $M\ w$) \iff
poincare-between $u\ v\ w$
proof (*cases* $u = v \vee v = w$)
case *True*
thus *?thesis*
using *assms*
unfolding *poincare-between-def*
by *auto*
next
case *False*
moreover
hence *moebius-pt* $M\ u \neq \text{moebius-pt } M\ v \wedge \text{moebius-pt } M\ v \neq \text{moebius-pt } M\ w$
by *auto*
moreover
have $v \neq \text{inversion } v\ w \neq \text{inversion } v$
using *inversion-noteq-unit-disc*[*of* $v\ w$]
using *inversion-noteq-unit-disc*[*of* $v\ v$]
using $\langle v \in \text{unit-disc} \rangle \langle w \in \text{unit-disc} \rangle$
by *auto*
ultimately
show *?thesis*
using *assms*
using *unit-circle-fix-moebius-pt-inversion*[*of* $M\ v$, *symmetric*]
unfolding *poincare-between-def*
by (*simp del: unit-circle-fix-moebius-pt-inversion*)
qed

lemma *conjugate-preserve-poincare-between* [*simp*]:
assumes $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$ **and** $w \in \text{unit-disc}$
shows *poincare-between* (*conjugate* u) (*conjugate* v) (*conjugate* w) \iff
poincare-between $u\ v\ w$
proof (*cases* $u = v \vee v = w$)
case *True*
thus *?thesis*
using *assms*
unfolding *poincare-between-def*
by *auto*
next
case *False*
moreover
hence *conjugate* $u \neq \text{conjugate } v \wedge \text{conjugate } v \neq \text{conjugate } w$
using *conjugate-inj* **by** *blast*
moreover
have $v \neq \text{inversion } v\ w \neq \text{inversion } v$
using *inversion-noteq-unit-disc*[*of* $v\ w$]
using *inversion-noteq-unit-disc*[*of* $v\ v$]
using $\langle v \in \text{unit-disc} \rangle \langle w \in \text{unit-disc} \rangle$

```

  by auto
ultimately
show ?thesis
  using assms
  using conjugate-cross-ratio[of v w inversion v u]
  unfolding poincare-between-def
  by (metis conjugate-id-iff conjugate-involution inversion-def inversion-sym o-apply)
qed

```

7.1.2 Some elementary properties of h-betweenness

```

lemma poincare-between-nonstrict [simp]:
  shows poincare-between u u v and poincare-between u v v
  by (simp-all add: poincare-between-def)

```

```

lemma poincare-between-sandwich:
  assumes u ∈ unit-disc and v ∈ unit-disc
  assumes poincare-between u v u
  shows u = v
proof (rule ccontr)
  assume ¬ ?thesis
  thus False
    using assms
    using inversion-noteq-unit-disc[of v u]
    using cross-ratio-1[of v u inversion v]
    unfolding poincare-between-def Let-def
    by auto
qed

```

```

lemma poincare-between-rev:
  assumes u ∈ unit-disc and v ∈ unit-disc and w ∈ unit-disc
  shows poincare-between u v w ⟷ poincare-between w v u
  using assms
  using inversion-noteq-unit-disc[of v w]
  using inversion-noteq-unit-disc[of v u]
  using cross-ratio-commute-13[of u v w inversion v]
  using cross-ratio-not-inf[of w inversion v v u]
  using cross-ratio-not-zero[of w v u inversion v]
  using inf-or-of-complex[of cross-ratio w v u (inversion v)]
  unfolding poincare-between-def
  by (auto simp add: Let-def Im-complex-div-eq-0 Re-divide divide-less-0-iff)

```

7.1.3 H-betweenness and h-collinearity

Three points can be in an h-between relation only when they are h-collinear.

```

lemma poincare-between-poincare-collinear [simp]:
  assumes in-disc: u ∈ unit-disc v ∈ unit-disc w ∈ unit-disc
  assumes betw: poincare-between u v w
  shows poincare-collinear {u, v, w}
proof (cases u = v ∨ v = w)
  case True
  thus ?thesis
    using assms
    by auto
next
  case False
  hence distinct: distinct [u, v, w, inversion v]
    using in-disc inversion-noteq-unit-disc[of v v] inversion-noteq-unit-disc[of v u] inversion-noteq-unit-disc[of v w]
    using betw poincare-between-sandwich[of w v]
    by (auto simp add: poincare-between-def Let-def)

then obtain H where *: {u, v, w, inversion v} ⊆ circline-set H
  using assms
  unfolding poincare-between-def
  using four-points-on-circline-iff-cross-ratio-real[of u v w inversion v]
  by auto

```

```

hence  $H = \text{poincare-line } u \ v$ 
  using assms distinct
  using unique-circline-set[of  $u \ v$  inversion  $v$ ]
  using poincare-line[of  $u \ v$ ] poincare-line-inversion[of  $u \ v$ ]
  unfolding circline-set-def
  by auto
thus ?thesis
  using * assms False
  unfolding poincare-collinear-def
  by (rule-tac  $x = \text{poincare-line } u \ v$  in exI) simp
qed

```

```

lemma poincare-between-poincare-line-uwz:
  assumes  $u \neq v$  and  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$  and
     $z \in \text{unit-disc}$  and poincare-between  $u \ v \ z$ 
  shows  $z \in \text{circline-set } (\text{poincare-line } u \ v)$ 
  using assms
  using poincare-between-poincare-collinear[of  $u \ v \ z$ ]
  using unique-poincare-line[OF assms(1-3)]
  unfolding poincare-collinear-def
  by auto

```

```

lemma poincare-between-poincare-line-uzv:
  assumes  $u \neq v$  and  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$  and
     $z \in \text{unit-disc}$  poincare-between  $u \ z \ v$ 
  shows  $z \in \text{circline-set } (\text{poincare-line } u \ v)$ 
  using assms
  using poincare-between-poincare-collinear[of  $u \ z \ v$ ]
  using unique-poincare-line[OF assms(1-3)]
  unfolding poincare-collinear-def
  by auto

```

7.1.4 H-betweenness on Euclidean segments

If the three points lie on an h-line that is a Euclidean line (e.g., if it contains zero), h-betweenness can be characterized much simpler than in the definition.

```

lemma poincare-between-x-axis-u0v:
  assumes is-real  $u'$  and  $u' \neq 0$  and  $v' \neq 0$ 
  shows poincare-between (of-complex  $u'$ )  $0_h$  (of-complex  $v'$ )  $\longleftrightarrow$  is-real  $v' \wedge \text{Re } u' * \text{Re } v' < 0$ 

```

proof –

```

  have  $\text{Re } u' \neq 0$ 
    using  $\langle \text{is-real } u' \rangle \langle u' \neq 0 \rangle$ 
    using complex-eq-if-Re-eq
    by auto
  have nz:  $\text{of-complex } u' \neq 0_h$  of-complex  $v' \neq 0_h$ 
    by (simp-all add:  $\langle u' \neq 0 \rangle \langle v' \neq 0 \rangle$ )
  hence  $0_h \neq \text{of-complex } v'$ 
    by metis

```

```

  let ?cr = cross-ratio (of-complex  $u'$ )  $0_h$  (of-complex  $v'$ )  $\infty_h$ 
  have is-real (to-complex ?cr)  $\wedge \text{Re } (\text{to-complex } ?cr) < 0 \longleftrightarrow$  is-real  $v' \wedge \text{Re } u' * \text{Re } v' < 0$ 
    using cross-ratio-0inf[of  $v' \ u'$ ]  $\langle v' \neq 0 \rangle \langle u' \neq 0 \rangle \langle \text{is-real } u' \rangle$ 
    by (metis Re-complex-div-lt-0 Re-mult-real complex-cnj-divide divide-cancel-left eq-cnj-iff-real to-complex-of-complex)
  thus ?thesis
    unfolding poincare-between-def inversion-zero
    using  $\langle \text{of-complex } u' \neq 0_h \rangle \langle 0_h \neq \text{of-complex } v' \rangle$ 
    by simp

```

qed

```

lemma poincare-between-u0v:
  assumes  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$  and  $u \neq 0_h$  and  $v \neq 0_h$ 
  shows poincare-between  $u \ 0_h \ v \longleftrightarrow$   $(\exists k < 0. \text{to-complex } u = \text{cor } k * \text{to-complex } v)$  (is  $?P \ u \ v$ )

```

proof (*cases* $u = v$)

```

  case True
  thus ?thesis
    using assms

```

```

using inf-or-of-complex[of v]
using poincare-between-sandwich[of u 0h]
by auto
next
case False
have  $\forall u. u \in \text{unit-disc} \wedge u \neq 0_h \longrightarrow ?P u v$  (is ?P' v)
proof (rule wlog-rotation-to-positive-x-axis)
  fix  $\varphi v$ 
  let ?M = moebius-pt (moebius-rotation  $\varphi$ )
  assume 1:  $v \in \text{unit-disc} v \neq 0_h$ 
  assume 2: ?P' (?M v)
  show ?P' v
  proof (rule allI, rule impI, (erule conjE)+)
    fix u
    assume 3:  $u \in \text{unit-disc} u \neq 0_h$ 
    have poincare-between (?M u) 0h (?M v)  $\longleftrightarrow$  poincare-between u 0h v
      using  $\langle u \in \text{unit-disc} \rangle \langle v \in \text{unit-disc} \rangle$ 
      using unit-disc-fix-moebius-preserve-poincare-between unit-disc-fix-rotation zero-in-unit-disc
      by fastforce
    thus ?P u v
      using 1 2 [rule-format, of ?M u] 3
      using inf-or-of-complex[of u] inf-or-of-complex[of v]
      by auto
  qed
next
fix x
assume 1: is-real x 0 < Re x Re x < 1
hence x  $\neq$  0
  by auto
show ?P' (of-complex x)
proof (rule allI, rule impI, (erule conjE)+)
  fix u
  assume 2:  $u \in \text{unit-disc} u \neq 0_h$ 
  then obtain u' where u = of-complex u'
    using inf-or-of-complex[of u]
    by auto
  show ?P u (of-complex x)
    using 1 2  $\langle x \neq 0 \rangle \langle u = \text{of-complex } u' \rangle$ 
    using poincare-between-rev[of u 0h of-complex x]
    using poincare-between-x-axis-u0v[of x u']  $\langle \text{is-real } x \rangle$ 
    apply (auto simp add: cmod-eq-Re)
    apply (rule-tac x=Re u' / Re x in exI, simp add: divide-neg-pos algebra-split-simps)
    using mult-neg-pos mult-pos-neg
    by blast
  qed
qed fact+
thus ?thesis
  using assms
  by auto
qed

lemma poincare-between-u0v-polar-form:
  assumes  $x \in \text{unit-disc}$  and  $y \in \text{unit-disc}$  and  $x \neq 0_h$  and  $y \neq 0_h$  and
    to-complex x = cor rx * cis  $\varphi$  to-complex y = cor ry * cis  $\varphi$ 
  shows poincare-between x 0h y  $\longleftrightarrow$  rx * ry < 0 (is ?P x y rx ry)
proof -
  from assms have rx  $\neq$  0 ry  $\neq$  0
    using inf-or-of-complex[of x] inf-or-of-complex[of y]
    by auto

  have  $(\exists k < 0. \text{cor } rx = \text{cor } k * \text{cor } ry) = (rx * ry < 0)$ 
  proof
    assume  $\exists k < 0. \text{cor } rx = \text{cor } k * \text{cor } ry$ 
    then obtain k where k < 0 cor rx = cor k * cor ry
      by auto
    hence rx = k * ry

```



```

    using of-real-eq-iff
    by fastforce
  thus  $rx * ry < 0$ 
    using  $\langle k < 0 \rangle \langle rx \neq 0 \rangle \langle ry \neq 0 \rangle$ 
    by (smt divisors-zero mult-nonneg-nonpos mult-nonpos-nonpos zero-less-mult-pos2)
next
  assume  $rx * ry < 0$ 
  hence  $rx = (rx/ry)*ry$   $rx / ry < 0$ 
    using  $\langle rx \neq 0 \rangle \langle ry \neq 0 \rangle$ 
    by (auto simp add: divide-less-0-iff algebra-split-simps)
  thus  $\exists k < 0. cor\ rx = cor\ k * cor\ ry$ 
    using  $\langle rx \neq 0 \rangle \langle ry \neq 0 \rangle$ 
    by (rule-tac  $x=rx / ry$  in  $exI, simp$ )
qed
  thus ?thesis
    using assms
    using poincare-between-u0v[OF assms(1-4)]
    by auto
qed

lemma poincare-between-x-axis-0uv:
  fixes  $x\ y :: real$ 
  assumes  $-1 < x$  and  $x < 1$  and  $x \neq 0$ 
  assumes  $-1 < y$  and  $y < 1$  and  $y \neq 0$ 
  shows poincare-between  $0_h$  (of-complex  $x$ ) (of-complex  $y$ )  $\longleftrightarrow$ 
    ( $x < 0 \wedge y < 0 \wedge y \leq x$ )  $\vee$  ( $x > 0 \wedge y > 0 \wedge x \leq y$ ) (is ?lhs  $\longleftrightarrow$  ?rhs)
proof (cases  $x = y$ )
  case True
  thus ?thesis
    using assms
    unfolding poincare-between-def
    by auto
next
  case False
  let ?x = of-complex  $x$  and ?y = of-complex  $y$ 

  have ?x  $\in$  unit-disc ?y  $\in$  unit-disc
    using assms
    by auto

  have distinct: distinct [ $0_h, ?x, ?y, inversion\ ?x$ ]
    using  $\langle x \neq 0 \rangle \langle y \neq 0 \rangle \langle x \neq y \rangle \langle ?x \in unit-disc \rangle \langle ?y \in unit-disc \rangle$ 
    using inversion-noteq-unit-disc[of ?x ?y]
    using inversion-noteq-unit-disc[of ?x ?x]
    using inversion-noteq-unit-disc[of ?x  $0_h$ ]
    using of-complex-inj[of  $x\ y$ ]
    by (metis distinct-length-2-or-more distinct-singleton of-complex-zero-iff of-real-eq-0-iff of-real-eq-iff zero-in-unit-disc)

  let ?cr = cross-ratio  $0_h$  ?x ?y (inversion ?x)
  have Re (to-complex ?cr) =  $x^2 * (x*y - 1) / (x * (y - x))$ 
    using  $\langle x \neq 0 \rangle \langle x \neq y \rangle$ 
    unfolding inversion-def
    by simp (transfer, transfer, auto simp add: vec-cnj-def power2-eq-square field-simps split: if-split-asm)
  moreover
  {
    fix  $a\ b :: real$ 
    assume  $b \neq 0$ 
    hence  $a < 0 \longleftrightarrow b^2 * a < (0::real)$ 
    by (metis mult.commute mult-eq-0-iff mult-neg-pos mult-pos-pos not-less-iff-gr-or-eq not-real-square-gt-zero power2-eq-square)
  }
  hence  $x^2 * (x*y - 1) < 0$ 
    using assms
    by (smt minus-mult-minus mult-le-cancel-left1)
  moreover
  have  $x * (y - x) > 0 \longleftrightarrow ?rhs$ 
    using  $\langle x \neq 0 \rangle \langle y \neq 0 \rangle \langle x \neq y \rangle$ 

```

```

    by (smt mult-le-0-iff)
ultimately
have *: Re (to-complex ?cr) < 0  $\longleftrightarrow$  ?rhs
    by (simp add: divide-less-0-iff)

show ?thesis
proof
  assume ?lhs
  have is-real (to-complex ?cr) Re (to-complex ?cr) < 0
    using <?lhs> distinct
    unfolding poincare-between-def Let-def
    by auto
  thus ?rhs
    using *
    by simp
next
  assume ?rhs
  hence Re (to-complex ?cr) < 0
    using *
    by simp
  moreover
  have {0h, of-complex (cor x), of-complex (cor y), inversion (of-complex (cor x))}  $\subseteq$  circline-set x-axis
    using <x  $\neq$  0> is-real-inversion[of cor x]
    using inf-or-of-complex[of inversion ?x]
    by (auto simp del: inversion-of-complex)
  hence is-real (to-complex ?cr)
    using four-points-on-circline-iff-cross-ratio-real[OF distinct]
    by auto
  ultimately
  show ?lhs
    using distinct
    unfolding poincare-between-def Let-def
    by auto
qed
qed

lemma poincare-between-0uv:
  assumes u  $\in$  unit-disc and v  $\in$  unit-disc and u  $\neq$  0h and v  $\neq$  0h
  shows poincare-between 0h u v  $\longleftrightarrow$ 
    (let u' = to-complex u; v' = to-complex v in Arg u' = Arg v'  $\wedge$  cmod u'  $\leq$  cmod v') (is ?P u v)
proof (cases u = v)
  case True
  thus ?thesis
    by simp
next
  case False
  have  $\forall v. v \in \text{unit-disc} \wedge v \neq 0_h \wedge v \neq u \longrightarrow (\text{poincare-between } 0_h \ u \ v \longleftrightarrow (\text{let } u' = \text{to-complex } u; v' = \text{to-complex } v \text{ in } \text{Arg } u' = \text{Arg } v' \wedge \text{cmod } u' \leq \text{cmod } v'))$  (is ?P' u)
  proof (rule wlog-rotation-to-positive-x-axis)
    show u  $\in$  unit-disc u  $\neq$  0h
      by fact+
  next
    fix x
    assume *: is-real x 0 < Re x Re x < 1
    hence of-complex x  $\in$  unit-disc of-complex x  $\neq$  0h of-complex x  $\in$  circline-set x-axis
      unfolding circline-set-x-axis
      by (auto simp add: cmod-eq-Re)
    show ?P' (of-complex x)
  proof safe
    fix v
    assume v  $\in$  unit-disc v  $\neq$  0h v  $\neq$  of-complex x poincare-between 0h (of-complex x) v
    hence v  $\in$  circline-set x-axis
      using poincare-between-poincare-line-uvz[of 0h of-complex x v]
      using poincare-line-0-real-is-x-axis[of of-complex x]
      using <of-complex x  $\neq$  0h> <v  $\neq$  0h> <v  $\neq$  of-complex x> <of-complex x  $\in$  unit-disc> <of-complex x  $\in$  circline-set x-axis>

```

```

    by auto
  obtain v' where v = of-complex v'
    using ⟨v ∈ unit-disc⟩
    using inf-or-of-complex[of v]
    by auto
  hence **: v = of-complex v' - 1 < Re v' Re v' < 1 Re v' ≠ 0 is-real v'
    using ⟨v ∈ unit-disc⟩ ⟨v ≠ 0h⟩ ⟨v ∈ circline-set x-axis⟩ of-complex-inj[of v']
    unfolding circline-set-x-axis
    by (auto simp add: cmod-eq-Re real-imag-0)
  show let u' = to-complex (of-complex x); v' = to-complex v in Arg u' = Arg v' ∧ cmod u' ≤ cmod v'
    using poincare-between-x-axis-0uv[of Re x Re v'] * **
    using ⟨poincare-between 0h (of-complex x) v⟩
    using arg-complex-of-real-positive[of Re x] arg-complex-of-real-negative[of Re x]
    using arg-complex-of-real-positive[of Re v'] arg-complex-of-real-negative[of Re v']
    by (auto simp add: cmod-eq-Re)
next
  fix v
  assume v ∈ unit-disc v ≠ 0h v ≠ of-complex x
  then obtain v' where **: v = of-complex v' v' ≠ 0 v' ≠ x
    using inf-or-of-complex[of v]
    by auto blast
  assume let u' = to-complex (of-complex x); v' = to-complex v in Arg u' = Arg v' ∧ cmod u' ≤ cmod v'
  hence ***: Re x < 0 ∧ Re v' < 0 ∧ Re v' ≤ Re x ∨ 0 < Re x ∧ 0 < Re v' ∧ Re x ≤ Re v' is-real v'
    using arg-pi-iff[of x] arg-pi-iff[of v']
    using arg-0-iff[of x] arg-0-iff[of v']
    using * **
    by (smt cmod-Re-le-iff to-complex-of-complex)+
  have -1 < Re v' Re v' < 1 Re v' ≠ 0 is-real v'
    using ⟨v ∈ unit-disc⟩ ** ⟨is-real v'⟩
    by (auto simp add: cmod-eq-Re complex-eq-if-Re-eq)
  thus poincare-between 0h (of-complex x) v
    using poincare-between-x-axis-0uv[of Re x Re v'] * ** ***
    by simp
qed
next
  fix φ u
  assume u ∈ unit-disc u ≠ 0h
  let ?M = moebius-rotation φ
  assume *: ?P' (moebius-pt ?M u)
  show ?P' u
  proof (rule allI, rule impI, (erule conjE)+)
    fix v
    assume v ∈ unit-disc v ≠ 0h v ≠ u
    have moebius-pt ?M v ≠ moebius-pt ?M u
      using ⟨v ≠ u⟩
      by auto
    obtain u' v' where v = of-complex v' u = of-complex u' v' ≠ 0 u' ≠ 0
      using inf-or-of-complex[of u] inf-or-of-complex[of v]
      using ⟨v ∈ unit-disc⟩ ⟨u ∈ unit-disc⟩ ⟨v ≠ 0h⟩ ⟨u ≠ 0h⟩
      by auto
    thus ?P u v
      using *[rule-format, of moebius-pt ?M v]
      using ⟨moebius-pt ?M v ≠ moebius-pt ?M u⟩
      using unit-disc-fix-moebius-preserve-poincare-between[of ?M 0h u v]
      using ⟨v ∈ unit-disc⟩ ⟨u ∈ unit-disc⟩ ⟨v ≠ 0h⟩ ⟨u ≠ 0h⟩
      using arg-mult-eq[of cis φ u' v']
      by simp (auto simp add: arg-mult norm-mult)
  qed
qed
  thus ?thesis
    using assms False
    by auto
qed
lemma poincare-between-y-axis-0uv:
  fixes x y :: real

```

```

assumes  $-1 < x$  and  $x < 1$  and  $x \neq 0$ 
assumes  $-1 < y$  and  $y < 1$  and  $y \neq 0$ 
shows poincare-between  $0_h$  (of-complex  $(i * x)$ ) (of-complex  $(i * y)$ )  $\longleftrightarrow$ 
  ( $x < 0 \wedge y < 0 \wedge y \leq x$ )  $\vee$  ( $x > 0 \wedge y > 0 \wedge x \leq y$ ) (is ?lhs  $\longleftrightarrow$  ?rhs)
using assms
using poincare-between-0uv[of of-complex  $(i * x)$  of-complex  $(i * y)$ ]
using arg-pi2-iff[of i * cor x] arg-pi2-iff[of i * cor y]
using arg-minus-pi2-iff[of i * cor x] arg-minus-pi2-iff[of i * cor y]
apply (simp add: norm-mult)
apply (smt (verit, best))
done

lemma poincare-between-x-axis-uvw:
fixes  $x y z :: \text{real}$ 
assumes  $-1 < x$  and  $x < 1$ 
assumes  $-1 < y$  and  $y < 1$  and  $y \neq x$ 
assumes  $-1 < z$  and  $z < 1$  and  $z \neq x$ 
shows poincare-between (of-complex  $x$ ) (of-complex  $y$ ) (of-complex  $z$ )  $\longleftrightarrow$ 
  ( $y < x \wedge z < x \wedge z \leq y$ )  $\vee$  ( $y > x \wedge z > x \wedge y \leq z$ ) (is ?lhs  $\longleftrightarrow$  ?rhs)
proof (cases  $x = 0 \vee y = 0 \vee z = 0$ )
case True
thus ?thesis
proof (cases  $x = 0$ )
case True
thus ?thesis
  using poincare-between-x-axis-0uv assms
  by simp
next
case False
show ?thesis
proof (cases  $z = 0$ )
case True
thus ?thesis
  using poincare-between-x-axis-0uv assms poincare-between-rev
  by (smt norm-of-real of-complex-zero of-real-0 poincare-between-nonstrict(2) unit-disc-iff-cmod-lt-1)
next
case False
have  $y = 0$ 
  using  $\langle x \neq 0 \rangle \langle z \neq 0 \rangle \langle x = 0 \vee y = 0 \vee z = 0 \rangle$ 
  by simp

have poincare-between (of-complex  $x$ )  $0_h$  (of-complex  $z$ ) = (is-real  $z \wedge x * z < 0$ )
  using  $\langle x \neq 0 \rangle \langle z \neq 0 \rangle$  poincare-between-x-axis-u0v
  by auto
moreover
have  $x * z < 0 \longleftrightarrow ?rhs$ 
  using True  $\langle x \neq 0 \rangle \langle z \neq 0 \rangle$ 
  by (smt zero-le-mult-iff)
ultimately
show ?thesis
  using  $\langle y = 0 \rangle$ 
  by auto
qed
qed
next
case False
thus ?thesis
proof (cases  $z = y$ )
case True
thus ?thesis
  using assms
  unfolding poincare-between-def
  by auto
next
case False
let ?x = of-complex  $x$  and ?y = of-complex  $y$  and ?z = of-complex  $z$ 

```

```

have ?x ∈ unit-disc ?y ∈ unit-disc ?z ∈ unit-disc
  using assms
  by auto

have distinct: distinct [?x, ?y, ?z, inversion ?y]
  using ⟨y ≠ x⟩ ⟨z ≠ x⟩ False ⟨?x ∈ unit-disc⟩ ⟨?y ∈ unit-disc⟩ ⟨?z ∈ unit-disc⟩
  using inversion-noteq-unit-disc[of ?y ?y]
  using inversion-noteq-unit-disc[of ?y ?x]
  using inversion-noteq-unit-disc[of ?y ?z]
  using of-complex-inj[of x y] of-complex-inj[of y z] of-complex-inj[of x z]
  by auto

have cor y * cor x ≠ 1
  using assms
  by (smt minus-mult-minus mult-less-cancel-left2 mult-less-cancel-right2 of-real-1 of-real-eq-iff of-real-mult)

let ?cr = cross-ratio ?x ?y ?z (inversion ?y)
have Re (to-complex ?cr) = (x - y) * (z*y - 1) / ((x*y - 1)*(z - y))
proof-
  have  $\bigwedge y x z. \llbracket y \neq x; z \neq x; z \neq y; \text{cor } y * \text{cor } x \neq 1; x \neq 0; y \neq 0; z \neq 0 \rrbracket \implies$ 
    (y * y + y * (y * (x * z)) - (y * x + y * (y * (y * z)))) /
    (y * y + y * (y * (x * z)) - (y * z + y * (y * (y * x)))) =
    (y + y * (x * z) - (x + y * (y * z))) / (y + y * (x * z) - (z + y * (y * x)))
  by (metis (no-types, opaque-lifting) ab-group-add-class.ab-diff-conv-add-uminus distrib-left mult-divide-mult-cancel-left-if
mult-minus-right)
  thus ?thesis
  using ⟨y ≠ x⟩ ⟨z ≠ x⟩ False ⟨¬ (x = 0 ∨ y = 0 ∨ z = 0)⟩
  using ⟨cor y * cor x ≠ 1⟩
  unfolding inversion-def
  by (transfer, transfer, auto simp add: vec-cnj-def power2-eq-square field-simps split: if-split-asm)
qed

moreover
have (x*y - 1) < 0
  using assms
  by (smt minus-mult-minus mult-less-cancel-right2 zero-less-mult-iff)
moreover
have (z*y - 1) < 0
  using assms
  by (smt minus-mult-minus mult-less-cancel-right2 zero-less-mult-iff)
moreover
have (x - y) / (z - y) < 0  $\longleftrightarrow$  ?rhs
  using ⟨y ≠ x⟩ ⟨z ≠ x⟩ False ⟨¬ (x = 0 ∨ y = 0 ∨ z = 0)⟩
  by (smt divide-less-cancel divide-nonneg-nonpos divide-nonneg-pos divide-nonpos-nonneg divide-nonpos-nonpos)
ultimately
have *: Re (to-complex ?cr) < 0  $\longleftrightarrow$  ?rhs
  by (smt algebra-split-simps(24) minus-divide-left zero-less-divide-iff zero-less-mult-iff)
show ?thesis
proof
  assume ?lhs
  have is-real (to-complex ?cr) Re (to-complex ?cr) < 0
    using ⟨?lhs⟩ distinct
    unfolding poincare-between-def Let-def
    by auto
  thus ?rhs
    using *
    by simp
next
  assume ?rhs
  hence Re (to-complex ?cr) < 0
    using *
    by simp
moreover
  have {of-complex (cor x), of-complex (cor y), of-complex (cor z), inversion (of-complex (cor y))}  $\subseteq$  circline-set
x-axis

```

```

    using ⟨ $\neg (x = 0 \vee y = 0 \vee z = 0)$ ⟩ is-real-inversion[of cor y]
    using inf-or-of-complex[of inversion ?y]
    by (auto simp del: inversion-of-complex)
  hence is-real (to-complex ?cr)
    using four-points-on-circline-iff-cross-ratio-real[OF distinct]
    by auto
  ultimately
  show ?lhs
    using distinct
    unfolding poincare-between-def Let-def
    by auto
qed
qed
qed

```

7.1.5 H-betweenness and h-collinearity

For three h-collinear points at least one of the three possible h-betweenness relations must hold.

lemma poincare-collinear3-between:

```

  assumes u ∈ unit-disc and v ∈ unit-disc and w ∈ unit-disc
  assumes poincare-collinear {u, v, w}
  shows poincare-between u v w ∨ poincare-between u w v ∨ poincare-between v u w (is ?P' u v w)
proof (cases u=v)
  case True
  thus ?thesis
    using assms
    by auto
next
  case False
  have ∀ w. w ∈ unit-disc ∧ poincare-collinear {u, v, w} ⟶ ?P' u v w (is ?P u v)
  proof (rule wlog-positive-x-axis[where P=?P])
    fix x
    assume x: is-real x 0 < Re x Re x < 1
    hence x ≠ 0
      using complex.expand[of x 0]
      by auto
    hence *: poincare-line 0h (of-complex x) = x-axis
      using x poincare-line-0-real-is-x-axis[of of-complex x]
      unfolding circline-set-x-axis
      by auto
    have of-complex x ∈ unit-disc
      using x
      by (auto simp add: cmod-eq-Re)
    have of-complex x ≠ 0h
      using ⟨x ≠ 0⟩
      by auto
    show ?P 0h (of-complex x)
  proof safe
    fix w
    assume w ∈ unit-disc
    assume poincare-collinear {0h, of-complex x, w}
    hence w ∈ circline-set x-axis
      using * unique-poincare-line[of 0h of-complex x] ⟨of-complex x ∈ unit-disc⟩ ⟨x ≠ 0⟩ ⟨of-complex x ≠ 0h⟩
      unfolding poincare-collinear-def
      by auto
    then obtain w' where w': w = of-complex w' is-real w'
      using ⟨w ∈ unit-disc⟩
      using inf-or-of-complex[of w]
      unfolding circline-set-x-axis
      by auto
    hence -1 < Re w' Re w' < 1
      using ⟨w ∈ unit-disc⟩
      by (auto simp add: cmod-eq-Re)
    assume 1: ¬ poincare-between (of-complex x) 0h w
    hence w ≠ 0h w' ≠ 0
      using w'

```

```

    unfolding poincare-between-def
  by auto
hence  $\operatorname{Re} w' \neq 0$ 
  using  $w'$  complex.expand[of  $w' 0$ ]
  by auto

have  $\operatorname{Re} w' \geq 0$ 
  using 1 poincare-between-x-axis-u0v[of  $x w'$ ]  $\langle \operatorname{Re} x > 0 \rangle$   $\langle \text{is-real } x \rangle$   $\langle x \neq 0 \rangle$   $\langle w' \neq 0 \rangle$   $w'$ 
  using mult-pos-neg
  by force

moreover

assume  $\neg$  poincare-between  $0_h$  (of-complex  $x$ )  $w$ 
hence  $\operatorname{Re} w' < \operatorname{Re} x$ 
  using poincare-between-x-axis-0uv[of  $\operatorname{Re} x \operatorname{Re} w'$ ]
  using  $w' x$   $\langle -1 < \operatorname{Re} w' \rangle$   $\langle \operatorname{Re} w' < 1 \rangle$   $\langle \operatorname{Re} w' \neq 0 \rangle$ 
  by auto

ultimately
show poincare-between  $0_h$   $w$  (of-complex  $x$ )
  using poincare-between-x-axis-0uv[of  $\operatorname{Re} w' \operatorname{Re} x$ ]
  using  $w' x$   $\langle -1 < \operatorname{Re} w' \rangle$   $\langle \operatorname{Re} w' < 1 \rangle$   $\langle \operatorname{Re} w' \neq 0 \rangle$ 
  by auto
qed
next
show  $u \in \text{unit-disc}$   $v \in \text{unit-disc}$   $u \neq v$ 
  by fact+
next
fix  $M u v$ 
assume 1:  $\text{unit-disc-fix } M u \in \text{unit-disc}$   $v \in \text{unit-disc}$   $u \neq v$ 
let  $?Mu = \text{moebius-pt } M u$  and  $?Mv = \text{moebius-pt } M v$ 
assume 2:  $?P ?Mu ?Mv$ 
show  $?P u v$ 
proof safe
  fix  $w$ 
  assume  $w \in \text{unit-disc}$   $\text{poincare-collinear } \{u, v, w\} \neg \text{poincare-between } u v w \neg \text{poincare-between } v u w$ 
  thus  $\text{poincare-between } u w v$ 
  using 1 2[rule-format, of  $\text{moebius-pt } M w$ ]
  by simp
qed
thus  $?thesis$ 
  using  $\text{assms}$ 
  by simp
qed

lemma poincare-collinear3-iff:
  assumes  $u \in \text{unit-disc}$   $v \in \text{unit-disc}$   $w \in \text{unit-disc}$ 
  shows  $\text{poincare-collinear } \{u, v, w\} \longleftrightarrow \text{poincare-between } u v w \vee \text{poincare-between } v u w \vee \text{poincare-between } v w u$ 
  using  $\text{assms}$ 
  by (metis  $\text{poincare-collinear3-between}$   $\text{insert-commute}$   $\text{poincare-between-poincare-collinear}$   $\text{poincare-between-rev}$ )

```

7.2 Some properties of betweenness

```

lemma poincare-between-transitivity:
  assumes  $a \in \text{unit-disc}$  and  $x \in \text{unit-disc}$  and  $b \in \text{unit-disc}$  and  $y \in \text{unit-disc}$  and
     $\text{poincare-between } a x b$  and  $\text{poincare-between } a b y$ 
  shows  $\text{poincare-between } x b y$ 
proof(cases  $a = b$ )
  case True
  thus  $?thesis$ 
    using  $\text{assms}$ 
    using  $\text{poincare-between-sandwich}$  by blast
next
  case False

```

```

have  $\forall x. \forall y. \text{poincare-between } a \ x \ b \wedge \text{poincare-between } a \ b \ y \wedge x \in \text{unit-disc}$ 
       $\wedge y \in \text{unit-disc} \longrightarrow \text{poincare-between } x \ b \ y$  (is ?P a b)
proof (rule wlog-positive-x-axis[where P=?P])
  show  $a \in \text{unit-disc}$ 
    using assms by simp
next
show  $b \in \text{unit-disc}$ 
  using assms by simp
next
show  $a \neq b$ 
  using False by simp
next
fix  $M \ u \ v$ 
assume *:  $\text{unit-disc-fix } M \ u \in \text{unit-disc} \ v \in \text{unit-disc} \ u \neq v$ 
       $\forall x \ y. \text{poincare-between } (\text{moebius-pt } M \ u) \ x \ (\text{moebius-pt } M \ v) \wedge$ 
       $\text{poincare-between } (\text{moebius-pt } M \ u) \ (\text{moebius-pt } M \ v) \ y \wedge$ 
       $x \in \text{unit-disc} \wedge y \in \text{unit-disc} \longrightarrow$ 
       $\text{poincare-between } x \ (\text{moebius-pt } M \ v) \ y$ 
show  $\forall x \ y. \text{poincare-between } u \ x \ v \wedge \text{poincare-between } u \ v \ y \wedge x \in \text{unit-disc} \wedge y \in \text{unit-disc}$ 
       $\longrightarrow \text{poincare-between } x \ v \ y$ 
proof safe
  fix  $x \ y$ 
  assume  $\text{poincare-between } u \ x \ v \ \text{poincare-between } u \ v \ y \ x \in \text{unit-disc} \ y \in \text{unit-disc}$ 

  have  $\text{poincare-between } (\text{moebius-pt } M \ u) \ (\text{moebius-pt } M \ x) \ (\text{moebius-pt } M \ v)$ 
    using  $\langle \text{poincare-between } u \ x \ v \rangle \langle \text{unit-disc-fix } M \rangle \langle x \in \text{unit-disc} \rangle \langle u \in \text{unit-disc} \rangle \langle v \in \text{unit-disc} \rangle$ 
    by simp
  moreover
  have  $\text{poincare-between } (\text{moebius-pt } M \ u) \ (\text{moebius-pt } M \ v) \ (\text{moebius-pt } M \ y)$ 
    using  $\langle \text{poincare-between } u \ v \ y \rangle \langle \text{unit-disc-fix } M \rangle \langle y \in \text{unit-disc} \rangle \langle u \in \text{unit-disc} \rangle \langle v \in \text{unit-disc} \rangle$ 
    by simp
  moreover
  have  $(\text{moebius-pt } M \ x) \in \text{unit-disc}$ 
    using  $\langle \text{unit-disc-fix } M \rangle \langle x \in \text{unit-disc} \rangle$  by simp
  moreover
  have  $(\text{moebius-pt } M \ y) \in \text{unit-disc}$ 
    using  $\langle \text{unit-disc-fix } M \rangle \langle y \in \text{unit-disc} \rangle$  by simp
  ultimately
  have  $\text{poincare-between } (\text{moebius-pt } M \ x) \ (\text{moebius-pt } M \ v) \ (\text{moebius-pt } M \ y)$ 
    using * by blast
  thus  $\text{poincare-between } x \ v \ y$ 
    using  $\langle y \in \text{unit-disc} \rangle * \langle x \in \text{unit-disc} \rangle$  by simp
qed
next
fix  $x$ 
assume  $xx: \text{is-real } x \ 0 < \text{Re } x \ \text{Re } x < 1$ 
hence  $\text{of-complex } x \in \text{unit-disc}$ 
  using cmod-eq-Re by auto
hence  $\text{of-complex } x \neq \infty_h$ 
  by simp
have  $\text{of-complex } x \neq 0_h$ 
  using xx by auto
have  $\text{of-complex } x \in \text{circline-set } x\text{-axis}$ 
  using xx by simp
show  $\forall m \ n. \text{poincare-between } 0_h \ m \ (\text{of-complex } x) \wedge \text{poincare-between } 0_h \ (\text{of-complex } x) \ n \wedge$ 
       $m \in \text{unit-disc} \wedge n \in \text{unit-disc} \longrightarrow \text{poincare-between } m \ (\text{of-complex } x) \ n$ 
proof safe
  fix  $m \ n$ 
  assume **:  $\text{poincare-between } 0_h \ m \ (\text{of-complex } x) \ \text{poincare-between } 0_h \ (\text{of-complex } x) \ n$ 
       $m \in \text{unit-disc} \ n \in \text{unit-disc}$ 
  show  $\text{poincare-between } m \ (\text{of-complex } x) \ n$ 
  proof(cases  $m = 0_h$ )
    case True
      thus ?thesis
        using ** by auto
  next

```



```

case False
hence  $m \in \text{circline-set } x\text{-axis}$ 
  using poincare-between-poincare-line-uzv[of  $0_h$  of-complex  $x$   $m$ ]
  using poincare-line-0-real-is-x-axis[of of-complex  $x$ ]
  using  $\langle \text{of-complex } x \in \text{unit-disc} \rangle \langle \text{of-complex } x \neq \infty_h \rangle \langle \text{of-complex } x \neq 0_h \rangle$ 
  using  $\langle \text{of-complex } x \in \text{circline-set } x\text{-axis} \rangle \langle m \in \text{unit-disc} \rangle ** (1)$ 
  by simp
then obtain  $m'$  where  $m = \text{of-complex } m'$  is-real  $m'$ 
  using inf-or-of-complex[of  $m$ ]  $\langle m \in \text{unit-disc} \rangle$ 
  unfolding circline-set-x-axis
  by auto
hence  $\text{Re } m' \leq \text{Re } x$ 
  using  $\langle \text{poincare-between } 0_h \text{ } m \text{ (of-complex } x) \rangle \langle \text{of-complex } x \neq 0_h \rangle$ 
  using False **  $\langle \text{of-complex } x \in \text{unit-disc} \rangle$ 
  using cmod-Re-le-iff poincare-between-0uv by auto

have  $n \neq 0_h$ 
  using  $** (2, 4) \langle \text{of-complex } x \neq 0_h \rangle \langle \text{of-complex } x \in \text{unit-disc} \rangle$ 
  using poincare-between-sandwich by fastforce
have  $n \in \text{circline-set } x\text{-axis}$ 
  using poincare-between-poincare-line-uvw[of  $0_h$  of-complex  $x$   $n$ ]
  using poincare-line-0-real-is-x-axis[of of-complex  $x$ ]
  using  $\langle \text{of-complex } x \in \text{unit-disc} \rangle \langle \text{of-complex } x \neq \infty_h \rangle \langle \text{of-complex } x \neq 0_h \rangle$ 
  using  $\langle \text{of-complex } x \in \text{circline-set } x\text{-axis} \rangle \langle n \in \text{unit-disc} \rangle ** (2)$ 
  by simp
then obtain  $n'$  where  $n = \text{of-complex } n'$  is-real  $n'$ 
  using inf-or-of-complex[of  $n$ ]  $\langle n \in \text{unit-disc} \rangle$ 
  unfolding circline-set-x-axis
  by auto
hence  $\text{Re } x \leq \text{Re } n'$ 
  using  $\langle \text{poincare-between } 0_h \text{ (of-complex } x) \text{ } n \rangle \langle \text{of-complex } x \neq 0_h \rangle$ 
  using False **  $\langle \text{of-complex } x \in \text{unit-disc} \rangle \langle n \neq 0_h \rangle$ 
  using cmod-Re-le-iff poincare-between-0uv
  by  $(\text{metis } \text{Re-complex-of-real } \text{arg-0-iff } \text{rcis-cmod-Arg } \text{rcis-zero-arg } \text{to-complex-of-complex})$ 

have poincare-between  $(\text{of-complex } m')$   $(\text{of-complex } x)$   $(\text{of-complex } n')$ 
  using  $\langle \text{Re } x \leq \text{Re } n' \rangle \langle \text{Re } m' \leq \text{Re } x \rangle$ 
  using poincare-between-x-axis-uvw[of  $\text{Re } m'$   $\text{Re } x$   $\text{Re } n'$ ]
  using  $\langle \text{is-real } n' \rangle \langle \text{is-real } m' \rangle \langle n \in \text{unit-disc} \rangle \langle n = \text{of-complex } n' \rangle$ 
  using  $\langle \text{of-complex } m' \rangle \langle m \in \text{unit-disc} \rangle$ 
  by  $(\text{smt } \text{complex-of-real-Re } \text{norm-of-real } \text{poincare-between-def } \text{unit-disc-iff-cmod-lt-1})$ 

thus ?thesis
  using  $\langle n = \text{of-complex } n' \rangle \langle m = \text{of-complex } m' \rangle$ 
  by auto
qed
qed
qed
thus ?thesis
  using assms
  by blast
qed

```

7.3 Poincare between - sum distances

Another possible definition of the h-betweenness relation is given in terms of h-distances between pairs of points. We prove it as a characterization equivalent to our cross-ratio based definition.

lemma *poincare-between-sum-distances-x-axis-u0v*:

assumes $\text{of-complex } u' \in \text{unit-disc}$ $\text{of-complex } v' \in \text{unit-disc}$

assumes $\text{is-real } u' \neq 0$ $v' \neq 0$

shows $\text{poincare-distance } (\text{of-complex } u') \text{ } 0_h + \text{poincare-distance } 0_h \text{ (of-complex } v') = \text{poincare-distance } (\text{of-complex } u') \text{ (of-complex } v') \iff$

$\text{is-real } v' \wedge \text{Re } u' * \text{Re } v' < 0$ (**is** $?P \text{ } u' \text{ } v' \iff ?Q \text{ } u' \text{ } v'$)

proof–

have $\text{Re } u' \neq 0$

```

using ‹is-real u'› ‹u' ≠ 0›
using complex-eq-if-Re-eq
by simp

let ?u = cmod u' and ?v = cmod v' and ?uv = cmod (u' - v')
have disc: ?u2 < 1 ?v2 < 1
  using unit-disc-cmod-square-lt-1[OF assms(1)]
  using unit-disc-cmod-square-lt-1[OF assms(2)]
  by auto
have poincare-distance (of-complex u') 0h + poincare-distance 0h (of-complex v') =
  arcosh (((1 + ?u2) * (1 + ?v2) + 4 * ?u * ?v) / ((1 - ?u2) * (1 - ?v2))) (is - = arcosh ?r1)
  using poincare-distance-formula-zero-sum[OF assms(1-2)]
  by (simp add: Let-def)
moreover
have poincare-distance (of-complex u') (of-complex v') =
  arcosh (((1 - ?u2) * (1 - ?v2) + 2 * ?uv2) / ((1 - ?u2) * (1 - ?v2))) (is - = arcosh ?r2)
  using disc
  using poincare-distance-formula[OF assms(1-2)]
  by (subst add-divide-distrib) simp
moreover
have arcosh ?r1 = arcosh ?r2 ‹⟷› ?Q u' v'
proof
  assume arcosh ?r1 = arcosh ?r2
  hence ?r1 = ?r2
  proof (subst (asm) arcosh-eq-iff)
    show ?r1 ≥ 1
    proof-
      have (1 - ?u2) * (1 - ?v2) ≤ (1 + ?u2) * (1 + ?v2) + 4 * ?u * ?v
      by (simp add: field-simps)
      thus ?thesis
      using disc
      by simp
    qed
  next
    show ?r2 ≥ 1
    using disc
    by simp
  qed
  hence (1 + ?u2) * (1 + ?v2) + 4 * ?u * ?v = (1 - ?u2) * (1 - ?v2) + 2 * ?uv2
  using disc
  by auto
  hence (cmod (u' - v'))2 = (cmod u' + cmod v')2
  by (simp add: field-simps power2-eq-square)
  hence *: Re u' * Re v' + |Re u'| * sqrt ((Im v')2 + (Re v')2) = 0
  using ‹is-real u'›
  unfolding cmod-power2 cmod-def
  by (simp add: field-simps) (simp add: power2-eq-square field-simps)
  hence sqrt ((Im v')2 + (Re v')2) = |Re v'|
  using ‹Re u' ≠ 0› ‹v' ≠ 0›
  by (smt complex-neq-0 mult.commute mult-cancel-right mult-minus-left real-sqrt-gt-0-iff)
  hence Im v' = 0
  by (smt Im-eq-0 norm-complex-def)
  moreover
  hence Re u' * Re v' = - |Re u'| * |Re v'|
  using *
  by simp
  hence Re u' * Re v' < 0
  using ‹Re u' ≠ 0› ‹v' ≠ 0›
  by (simp add: ‹is-real v'› complex-eq-if-Re-eq)
  ultimately
  show ?Q u' v'
  by simp
next
  assume ?Q u' v'
  hence is-real v' Re u' * Re v' < 0
  by auto

```

```

have ?r1 = ?r2
proof (cases Re u' > 0)
  case True
  hence Re v' < 0
  using ⟨Re u' * Re v' < 0⟩
  by (smt zero-le-mult-iff)
  show ?thesis
  using disc ⟨is-real u'⟩ ⟨is-real v'⟩
  using ⟨Re u' > 0⟩ ⟨Re v' < 0⟩
  unfolding cmod-power2 cmod-def
  by simp (simp add: power2-eq-square field-simps)
next
  case False
  hence Re u' < 0
  using ⟨Re u' ≠ 0⟩
  by simp
  hence Re v' > 0
  using ⟨Re u' * Re v' < 0⟩
  by (smt zero-le-mult-iff)
  show ?thesis
  using disc ⟨is-real u'⟩ ⟨is-real v'⟩
  using ⟨Re u' < 0⟩ ⟨Re v' > 0⟩
  unfolding cmod-power2 cmod-def
  by simp (simp add: power2-eq-square field-simps)
qed
thus arcosh ?r1 = arcosh ?r2
  by metis
qed
ultimately
show ?thesis
  by simp
qed

```

Different proof of the previous theorem relying on the cross-ratio definition, and not the distance formula. We suppose that this could be also used to prove the triangle inequality.

lemma *poincare-between-sum-distances-x-axis-u0v-different-proof*:

```

assumes of-complex u' ∈ unit-disc of-complex v' ∈ unit-disc
assumes is-real u' u' ≠ 0 v' ≠ 0 is-real v'
shows poincare-distance (of-complex u') 0h + poincare-distance 0h (of-complex v') = poincare-distance (of-complex u') (of-complex v') ⟷
  Re u' * Re v' < 0 (is ?P u' v' ⟷ ?Q u' v')

```

proof–

```

have -1 < Re u' Re u' < 1 Re u' ≠ 0
  using assms
  by (auto simp add: cmod-eq-Re complex-eq-if-Re-eq)
have -1 < Re v' Re v' < 1 Re v' ≠ 0
  using assms
  by (auto simp add: cmod-eq-Re complex-eq-if-Re-eq)

```

```

have |ln (Re ((1 - u') / (1 + u')))| + |ln (Re ((1 - v') / (1 + v')))| =
  |ln (Re ((1 + u') * (1 - v') / ((1 - u') * (1 + v')))| ⟷ Re u' * Re v' < 0 (is |ln ?a1| + |ln ?a2| = |ln ?a3|)
⟷ -)

```

proof–

```

have 1: 0 < ?a1 ln ?a1 > 0 ⟷ Re u' < 0
  using ⟨Re u' < 1⟩ ⟨Re u' > -1⟩ ⟨is-real u'⟩
  using complex-is-Real-iff
  by auto
have 2: 0 < ?a2 ln ?a2 > 0 ⟷ Re v' < 0
  using ⟨Re v' < 1⟩ ⟨Re v' > -1⟩ ⟨is-real v'⟩
  using complex-is-Real-iff
  by auto
have 3: 0 < ?a3 ln ?a3 > 0 ⟷ Re v' < Re u'
  using ⟨Re u' < 1⟩ ⟨Re u' > -1⟩ ⟨is-real u'⟩
  using ⟨Re v' < 1⟩ ⟨Re v' > -1⟩ ⟨is-real v'⟩
  using complex-is-Real-iff
  by auto (simp add: field-simps)+

```

```

show ?thesis
proof
  assume *:  $Re\ u' * Re\ v' < 0$ 
  show  $|\ln\ ?a1| + |\ln\ ?a2| = |\ln\ ?a3|$ 
  proof (cases  $Re\ u' > 0$ )
    case True
      hence  $Re\ v' < 0$ 
      using *
      by (smt mult-nonneg-nonneg)
    show ?thesis
      using 1 2 3  $\langle Re\ u' > 0 \rangle \langle Re\ v' < 0 \rangle$ 
      using  $\langle Re\ u' < 1 \rangle \langle Re\ u' > -1 \rangle \langle is-real\ u' \rangle$ 
      using  $\langle Re\ v' < 1 \rangle \langle Re\ v' > -1 \rangle \langle is-real\ v' \rangle$ 
      using complex-is-Real-iff
      using ln-div ln-mult
      by simp
  next
    case False
      hence  $Re\ v' > 0\ Re\ u' < 0$ 
      using *
      by (smt zero-le-mult-iff)+
    show ?thesis
      using 1 2 3  $\langle Re\ u' < 0 \rangle \langle Re\ v' > 0 \rangle$ 
      using  $\langle Re\ u' < 1 \rangle \langle Re\ u' > -1 \rangle \langle is-real\ u' \rangle$ 
      using  $\langle Re\ v' < 1 \rangle \langle Re\ v' > -1 \rangle \langle is-real\ v' \rangle$ 
      using complex-is-Real-iff
      using ln-div ln-mult
      by simp
  qed
next
assume *:  $|\ln\ ?a1| + |\ln\ ?a2| = |\ln\ ?a3|$ 
{
  assume  $Re\ u' > 0\ Re\ v' > 0$ 
  hence False
  using * 1 2 3
  using  $\langle Re\ u' < 1 \rangle \langle Re\ u' > -1 \rangle \langle is-real\ u' \rangle$ 
  using  $\langle Re\ v' < 1 \rangle \langle Re\ v' > -1 \rangle \langle is-real\ v' \rangle$ 
  using complex-is-Real-iff
  using ln-mult ln-div
  by (cases  $Re\ v' < Re\ u'$ ) auto
}
}
moreover
{
  assume  $Re\ u' < 0\ Re\ v' < 0$ 
  hence False
  using * 1 2 3
  using  $\langle Re\ u' < 1 \rangle \langle Re\ u' > -1 \rangle \langle is-real\ u' \rangle$ 
  using  $\langle Re\ v' < 1 \rangle \langle Re\ v' > -1 \rangle \langle is-real\ v' \rangle$ 
  using complex-is-Real-iff
  using ln-mult ln-div
  by (cases  $Re\ v' < Re\ u'$ ) auto
}
}
ultimately
show  $Re\ u' * Re\ v' < 0$ 
  using  $\langle Re\ u' \neq 0 \rangle \langle Re\ v' \neq 0 \rangle$ 
  by (smt divisors-zero mult-le-0-iff)
qed
qed
thus ?thesis
  using assms
  apply (subst poincare-distance-sym, simp, simp)
  apply (subst poincare-distance-zero-x-axis, simp, simp add: circline-set-x-axis)
  apply (subst poincare-distance-zero-x-axis, simp, simp add: circline-set-x-axis)
  apply (subst poincare-distance-x-axis-x-axis, simp, simp, simp add: circline-set-x-axis, simp add: circline-set-x-axis)
  apply simp
done

```

qed

lemma *poincare-between-sum-distances*:

assumes $u \in \text{unit-disc}$ and $v \in \text{unit-disc}$ and $w \in \text{unit-disc}$

shows *poincare-between* $u v w \longleftrightarrow$

$\text{poincare-distance } u v + \text{poincare-distance } v w = \text{poincare-distance } u w$ (is $?P' u v w$)

proof (cases $u = v$)

case *True*

thus *?thesis*

using *assms*

by *simp*

next

case *False*

have $\forall w. w \in \text{unit-disc} \longrightarrow (\text{poincare-between } u v w \longleftrightarrow \text{poincare-distance } u v + \text{poincare-distance } v w = \text{poincare-distance } u w)$ (is $?P u v$)

proof (rule *wlog-positive-x-axis*)

fix x

assume *is-real* x $0 < \text{Re } x$ $\text{Re } x < 1$

have *of-complex* $x \in \text{circline-set } x\text{-axis}$

using $\langle \text{is-real } x \rangle$

by (*auto simp add: circline-set-x-axis*)

have *of-complex* $x \in \text{unit-disc}$

using $\langle \text{is-real } x \rangle \langle 0 < \text{Re } x \rangle \langle \text{Re } x < 1 \rangle$

by (*simp add: cmod-eq-Re*)

have $x \neq 0$

using $\langle \text{is-real } x \rangle \langle \text{Re } x > 0 \rangle$

by *auto*

show $?P$ (*of-complex* x) 0_h

proof (rule *allI*, rule *impI*)

fix w

assume $w \in \text{unit-disc}$

then obtain w' where $w = \text{of-complex } w'$

using *inf-or-of-complex*[*of* w]

by *auto*

show $?P'$ (*of-complex* x) $0_h w$

proof (cases $w = 0_h$)

case *True*

thus *?thesis*

by *simp*

next

case *False*

hence $w' \neq 0$

using $\langle w = \text{of-complex } w' \rangle$

by *auto*

show *?thesis*

using $\langle \text{is-real } x \rangle \langle x \neq 0 \rangle \langle w = \text{of-complex } w' \rangle \langle w' \neq 0 \rangle$

using $\langle \text{of-complex } x \in \text{unit-disc} \rangle \langle w \in \text{unit-disc} \rangle$

apply *simp*

apply (*subst poincare-between-x-axis-u0v*, *simp-all*)

apply (*subst poincare-between-sum-distances-x-axis-u0v*, *simp-all*)

done

qed

qed

next

show $v \in \text{unit-disc}$ $u \in \text{unit-disc}$

using *assms*

by *auto*

next

show $v \neq u$

using $\langle u \neq v \rangle$

by *simp*

```

next
  fix M u v
  assume *: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc u ≠ v and
    **: ?P (moebius-pt M v) (moebius-pt M u)
  show ?P v u
  proof (rule allI, rule impI)
    fix w
    assume w ∈ unit-disc
    hence moebius-pt M w ∈ unit-disc
      using ⟨unit-disc-fix M⟩
      by auto
    thus ?P' v u w
      using ⟨u ∈ unit-disc⟩ ⟨v ∈ unit-disc⟩ ⟨w ∈ unit-disc⟩ ⟨unit-disc-fix M⟩
      using **[rule-format, of moebius-pt M w]
      by auto
  qed
qed
thus ?thesis
  using assms
  by simp
qed

```

7.4 Some more properties of h-betweenness.

Some lemmas proved earlier are proved almost directly using the sum of distances characterization.

```

lemma unit-disc-fix-moebius-preserve-poincare-between':
  assumes unit-disc-fix M and u ∈ unit-disc and v ∈ unit-disc and w ∈ unit-disc
  shows poincare-between (moebius-pt M u) (moebius-pt M v) (moebius-pt M w) ⟷
    poincare-between u v w
  using assms
  using poincare-between-sum-distances
  by simp

```

```

lemma conjugate-preserve-poincare-between':
  assumes u ∈ unit-disc v ∈ unit-disc w ∈ unit-disc
  shows poincare-between (conjugate u) (conjugate v) (conjugate w) ⟷ poincare-between u v w
  using assms
  using poincare-between-sum-distances
  by simp

```

There is a unique point on a ray on the given distance from the given starting point

```

lemma unique-poincare-distance-on-ray:
  assumes d ≥ 0 u ≠ v u ∈ unit-disc v ∈ unit-disc
  assumes y ∈ unit-disc poincare-distance u y = d poincare-between u v y
  assumes z ∈ unit-disc poincare-distance u z = d poincare-between u v z
  shows y = z
proof -
  have ∀ d y z. d ≥ 0 ∧
    y ∈ unit-disc ∧ poincare-distance u y = d ∧ poincare-between u v y ∧
    z ∈ unit-disc ∧ poincare-distance u z = d ∧ poincare-between u v z ⟶ y = z (is ?P u v)
proof (rule wlog-positive-x-axis[where P=?P])
  fix x
  assume x: is-real x 0 < Re x Re x < 1
  hence x ≠ 0
    using complex.expand[of x 0]
    by auto
  hence *: poincare-line 0h (of-complex x) = x-axis
    using x poincare-line-0-real-is-x-axis[of of-complex x]
    unfolding circline-set-x-axis
    by auto
  have of-complex x ∈ unit-disc
    using x
    by (auto simp add: cmod-eq-Re)
  have Arg x = 0
    using x

```

```

    using arg-0-iff by blast
  show ?P 0h (of-complex x)
  proof safe
    fix y z
    assume y ∈ unit-disc z ∈ unit-disc
    then obtain y' z' where yz: y = of-complex y' z = of-complex z'
      using inf-or-of-complex[of y] inf-or-of-complex[of z]
      by auto
    assume betw: poincare-between 0h (of-complex x) y poincare-between 0h (of-complex x) z
    hence y ≠ 0h z ≠ 0h
      using ⟨x ≠ 0⟩ ⟨of-complex x ∈ unit-disc⟩ ⟨y ∈ unit-disc⟩
      using poincare-between-sandwich[of 0h of-complex x]
      using of-complex-zero-iff[of x]
      by force+

    hence Arg y' = 0 cmod y' ≥ cmod x Arg z' = 0 cmod z' ≥ cmod x
      using poincare-between-0uw[of of-complex x y] poincare-between-0uw[of of-complex x z]
      using ⟨of-complex x ∈ unit-disc⟩ ⟨x ≠ 0⟩ ⟨Arg x = 0⟩ ⟨y ∈ unit-disc⟩ ⟨z ∈ unit-disc⟩ betw yz
      by (simp-all add: Let-def)
    hence *: is-real y' is-real z' Re y' > 0 Re z' > 0
      using arg-0-iff[of y'] arg-0-iff[of z'] x ⟨y ≠ 0h⟩ ⟨z ≠ 0h⟩ yz
      by auto
    assume poincare-distance 0h z = poincare-distance 0h y 0 ≤ poincare-distance 0h y
    thus y = z
      using * yz ⟨y ∈ unit-disc⟩ ⟨z ∈ unit-disc⟩
      using unique-x-axis-poincare-distance-positive[of poincare-distance 0h y]
      by (auto simp add: cmod-eq-Re unit-disc-to-complex-inj)
  qed
next
  show u ∈ unit-disc v ∈ unit-disc u ≠ v
    by fact+
next
  fix M u v
  assume *: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc u ≠ v
  assume **: ?P (moebius-pt M u) (moebius-pt M v)
  show ?P u v
  proof safe
    fix d y z
    assume ***: 0 ≤ poincare-distance u y
      y ∈ unit-disc poincare-between u v y
      z ∈ unit-disc poincare-between u v z
      poincare-distance u z = poincare-distance u y
    let ?Mu = moebius-pt M u and ?Mv = moebius-pt M v and ?My = moebius-pt M y and ?Mz = moebius-pt M z
    have ?Mu ∈ unit-disc ?Mv ∈ unit-disc ?My ∈ unit-disc ?Mz ∈ unit-disc
      using ⟨u ∈ unit-disc⟩ ⟨v ∈ unit-disc⟩ ⟨y ∈ unit-disc⟩ ⟨z ∈ unit-disc⟩
      using ⟨unit-disc-fix M⟩
      by auto
    hence ?My = ?Mz
      using * ***
      using **[rule-format, of poincare-distance ?Mu ?My ?My ?Mz]
      by simp
    thus y = z
      using bij-moebius-pt[of M]
      unfolding bij-def inj-on-def
      by blast
  qed
qed
thus ?thesis
  using assms
  by auto
qed

end
theory Poincare-Lines-Axis-Intersections
  imports Poincare-Between
begin

```

8 Intersection of h-lines with the x-axis in the Poincaré model

8.1 Betweenness of x-axis intersection

The intersection point of the h-line determined by points u and v and the x-axis is between u and v , then u and v are in the opposite half-planes (one must be in the upper, and the other one in the lower half-plane).

lemma *poincare-between-x-axis-intersection*:

assumes $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$ **and** $z \in \text{unit-disc}$ **and** $u \neq v$
assumes $u \notin \text{circline-set x-axis}$ **and** $v \notin \text{circline-set x-axis}$
assumes $z \in \text{circline-set (poincare-line } u \ v) \cap \text{circline-set x-axis}$
shows *poincare-between* $u \ z \ v \iff \text{Arg (to-complex } u) * \text{Arg (to-complex } v) < 0$

proof–

have $\forall u \ v. u \in \text{unit-disc} \wedge v \in \text{unit-disc} \wedge u \neq v \wedge$
 $u \notin \text{circline-set x-axis} \wedge v \notin \text{circline-set x-axis} \wedge$
 $z \in \text{circline-set (poincare-line } u \ v) \cap \text{circline-set x-axis} \longrightarrow$
 $(\text{poincare-between } u \ z \ v \iff \text{Arg (to-complex } u) * \text{Arg (to-complex } v) < 0)$ **(is ?P z)**

proof (rule *wlog-real-zero*)

show *?P 0_h*

proof ((rule *allI*)+, rule *impI*, (erule *conjE*)+)

fix $u \ v$

assume $*$: $u \in \text{unit-disc}$ $v \in \text{unit-disc}$ $u \neq v$
 $u \notin \text{circline-set x-axis}$ $v \notin \text{circline-set x-axis}$
 $0_h \in \text{circline-set (poincare-line } u \ v) \cap \text{circline-set x-axis}$

obtain $u' \ v'$ **where** $uw: u = \text{of-complex } u' \ v = \text{of-complex } v'$

using $*$ *inf-or-of-complex[of u]* *inf-or-of-complex[of v]*

by *auto*

hence $u \neq 0_h$ $v \neq 0_h$ $u' \neq 0$ $v' \neq 0$

using $*$

by *auto*

hence $\text{Arg } u' \neq 0$ $\text{Arg } v' \neq 0$

using $*$ *arg-0-iff[of u']* *arg-0-iff[of v']*

unfolding *circline-set-x-axis uw*

by *auto*

have *poincare-collinear* $\{0_h, u, v\}$

using $*$

unfolding *poincare-collinear-def*

by (rule-tac $x = \text{poincare-line } u \ v$ **in** *exI*, *simp*)

have $(\exists k < 0. u' = \text{cor } k * v') \iff (\text{Arg } u' * \text{Arg } v' < 0)$ **(is ?lhs \iff ?rhs)**

proof

assume *?lhs*

then obtain k **where** $k < 0$ $u' = \text{cor } k * v'$

by *auto*

thus *?rhs*

using *arg-mult-real-negative[of k v']* *arg-uminus-opposite-sign[of v']*

using $\langle u' \neq 0 \rangle$ $\langle v' \neq 0 \rangle$ $\langle \text{Arg } u' \neq 0 \rangle$ $\langle \text{Arg } v' \neq 0 \rangle$

by (auto *simp add: mult-neg-pos mult-pos-neg*)

next

assume *?rhs*

obtain $ru \ rv \ \varphi$ **where** *polar*: $u' = \text{cor } ru * \text{cis } \varphi$ $v' = \text{cor } rv * \text{cis } \varphi$

using $\langle \text{poincare-collinear } \{0_h, u, v\} \rangle$ *poincare-collinear-zero-polar-form[of u' v']* $uw * \langle u' \neq 0 \rangle \langle v' \neq 0 \rangle$

by *auto*

have $ru * rv < 0$

using *polar* $\langle ?rhs \rangle$ $\langle u' \neq 0 \rangle$ $\langle v' \neq 0 \rangle$

using *arg-mult-real-negative[of ru cis φ]* *arg-mult-real-positive[of ru cis φ]*

using *arg-mult-real-negative[of rv cis φ]* *arg-mult-real-positive[of rv cis φ]*

apply (cases $ru > 0$)

apply (cases $rv > 0$, *simp*, *simp add: mult-pos-neg*)

apply (cases $rv > 0$, *simp add: mult-neg-pos*, *simp*)

done

thus *?lhs*

using *polar*

by (rule-tac $x = ru / rv$ **in** *exI*, *auto simp add: divide-less-0-iff mult-less-0-iff*)


```

qed
thus poincare-between u 0h v = (Arg (to-complex u) * Arg (to-complex v) < 0)
  using poincare-between-u0v[of u v] * ⟨u ≠ 0h⟩ ⟨v ≠ 0h⟩ uv
  by simp
qed
next
fix a z
assume 1: is-real a cmod a < 1 z ∈ unit-disc
assume 2: ?P (moebius-pt (blaschke a) z)
show ?P z
proof ((rule allI)+, rule impI, (erule conjE)+)
  fix u v
  let ?M = moebius-pt (blaschke a)
  let ?Mu = ?M u
  let ?Mv = ?M v
  assume *: u ∈ unit-disc v ∈ unit-disc u ≠ v u ∉ circline-set x-axis v ∉ circline-set x-axis
  hence u ≠ ∞h v ≠ ∞h
  by auto

  have **: ∧ x y :: real. x * y < 0 ⟷ sgn (x * y) < 0
  by simp

  assume z ∈ circline-set (poincare-line u v) ∩ circline-set x-axis
  thus poincare-between u z v = (Arg (to-complex u) * Arg (to-complex v) < 0)
    using * 1 2[rule-format, of ?Mu ?Mv] ⟨cmod a < 1⟩ ⟨is-real a⟩ blaschke-unit-disc-fix[of a]
    using inversion-noteq-unit-disc[of of-complex a u] ⟨u ≠ ∞h⟩
    using inversion-noteq-unit-disc[of of-complex a v] ⟨v ≠ ∞h⟩
    apply auto
  apply (subst (asm) **, subst **, subst (asm) sgn-mult, subst sgn-mult, simp)
  apply (subst (asm) **, subst (asm) **, subst (asm) sgn-mult, subst (asm) sgn-mult, simp)
  done
qed
next
show z ∈ unit-disc by fact
next
show is-real (to-complex z)
  using assms inf-or-of-complex[of z]
  by (auto simp add: circline-set-x-axis)
qed
thus ?thesis
  using assms
  by simp
qed

```

8.2 Check if an h-line intersects the x-axis

lemma *x-axis-intersection-equation*:

```

assumes
  H = mk-circline A B C D and
  (A, B, C, D) ∈ hermitean-nonzero
shows of-complex z ∈ circline-set x-axis ∩ circline-set H ⟷
  A*z2 + 2*Re B*z + D = 0 ∧ is-real z (is ?lhs ⟷ ?rhs)
proof -
  have ?lhs ⟷ A*z2 + (B + cnj B)*z + D = 0 ∧ z = cnj z
  using assms
  using circline-equation-x-axis[of z]
  using circline-equation[of H A B C D z]
  using hermitean-elems
  by (auto simp add: power2-eq-square field-simps)
  thus ?thesis
  using eq-cnj-iff-real[of z]
  using hermitean-elems[of A B C D]
  by (simp add: complex-add-cnj complex-eq-if-Re-eq)
qed

```

Check if an h-line intersects x-axis within the unit disc - this could be generalized to checking if an arbitrary

circle intersects the x-axis, but we do not need that.

definition *intersects-x-axis-cmat* :: *complex-mat* \Rightarrow *bool* **where**
 [simp]: *intersects-x-axis-cmat* $H = (\text{let } (A, B, C, D) = H \text{ in } A = 0 \vee (\text{Re } B)^2 > (\text{Re } A)^2)$

lift-definition *intersects-x-axis-clmat* :: *circle-mat* \Rightarrow *bool* **is** *intersects-x-axis-cmat*
 done

lift-definition *intersects-x-axis* :: *circle* \Rightarrow *bool* **is** *intersects-x-axis-clmat*

proof (*transfer*)

fix $H1\ H2$

assume $hh: \text{hermitean } H1 \wedge H1 \neq \text{mat-zero}$ **and** $\text{hermitean } H2 \wedge H2 \neq \text{mat-zero}$

obtain $A1\ B1\ C1\ D1\ A2\ B2\ C2\ D2$ **where** $*$: $H1 = (A1, B1, C1, D1)\ H2 = (A2, B2, C2, D2)$

by (*cases* $H1$, *cases* $H2$, *auto*)

assume *circle-eq-cmat* $H1\ H2$

then obtain k **where** $k: k \neq 0 \wedge H2 = \text{cor } k *_{sm} H1$

by *auto*

show *intersects-x-axis-cmat* $H1 = \text{intersects-x-axis-cmat } H2$

proof–

have $k \neq 0 \implies (\text{Re } A1)^2 < (\text{Re } B1)^2 \iff (k * \text{Re } A1)^2 < (k * \text{Re } B1)^2$

by (*smt mult-strict-left-mono power2-eq-square semiring-normalization-rules(13) zero-less-power2*)

thus *?thesis*

using $*$ k

by *auto*

qed

qed

lemma *intersects-x-axis-mk-circle*:

assumes *is-real* A **and** $A \neq 0 \vee B \neq 0$

shows *intersects-x-axis* (*mk-circle* $A\ B\ (\text{cnj } B)\ A$) $\iff A = 0 \vee (\text{Re } B)^2 > (\text{Re } A)^2$

proof–

let $?H = (A, B, (\text{cnj } B), A)$

have *hermitean* $?H$

using $\langle \text{is-real } A \rangle$

unfolding *hermitean-def mat-adj-def mat-cnj-def*

using *eq-cnj-iff-real*

by *auto*

moreover

have $?H \neq \text{mat-zero}$

using *assms*

by *auto*

ultimately

show *?thesis*

by (*transfer, transfer, auto simp add: Let-def*)

qed

lemma *intersects-x-axis-iff*:

assumes *is-poincare-line* H

shows $(\exists x \in \text{unit-disc. } x \in \text{circle-set } H \cap \text{circle-set } x\text{-axis}) \iff \text{intersects-x-axis } H$

proof–

obtain $Ac\ B\ C\ Dc$ **where** $*$: $H = \text{mk-circle } Ac\ B\ C\ Dc$ *hermitean* (Ac, B, C, Dc) $(Ac, B, C, Dc) \neq \text{mat-zero}$

using *ex-mk-circle[of H]*

by *auto*

hence $(\text{cmod } B)^2 > (\text{cmod } Ac)^2 \wedge Ac = Dc$

using *assms*

using *is-poincare-line-mk-circle*

by *auto*

hence $H = \text{mk-circle } (\text{Re } Ac)\ B\ (\text{cnj } B)\ (\text{Re } Ac)$ *hermitean* $(\text{cor } (\text{Re } Ac), B, (\text{cnj } B), \text{cor } (\text{Re } Ac))$ $(\text{cor } (\text{Re } Ac), B, (\text{cnj } B), \text{cor } (\text{Re } Ac)) \neq \text{mat-zero}$

using *hermitean-elems[of Ac B C Dc]* $*$

by *auto*

then obtain A **where**

$*$: $H = \text{mk-circle } (\text{cor } A)\ B\ (\text{cnj } B)\ (\text{cor } A)$ $(\text{cor } A, B, (\text{cnj } B), \text{cor } A) \in \text{hermitean-nonzero}$

by *auto*

show *?thesis*

```

proof (cases A = 0)
  case True
  thus ?thesis
  using *
  using x-axis-intersection-equation[OF *(1-2), of 0]
  using intersects-x-axis-mk-circline[of cor A B]
  by auto
next
case False
show ?thesis
proof
  assume  $\exists x \in \text{unit-disc. } x \in \text{circline-set } H \cap \text{circline-set } x\text{-axis}$ 
  then obtain x where **: of-complex x  $\in$  unit-disc of-complex x  $\in$  circline-set H  $\cap$  circline-set x-axis
  by (metis inf-or-of-complex inf-notin-unit-disc)
  hence is-real x
  unfolding circline-set-x-axis
  using of-complex-inj
  by auto
  hence eq:  $A * (\text{Re } x)^2 + 2 * \text{Re } B * \text{Re } x + A = 0$ 
  using **
  using x-axis-intersection-equation[OF *(1-2), of Re x]
  by simp
  hence  $(2 * \text{Re } B)^2 - 4 * A * A \geq 0$ 
  using discriminant-iff[of A - 2 * Re B A]
  using discrim-def[of A 2 * Re B A] False
  by auto
  hence  $(\text{Re } B)^2 \geq (\text{Re } A)^2$ 
  by (simp add: power2-eq-square)
  moreover
  have  $(\text{Re } B)^2 \neq (\text{Re } A)^2$ 
  proof (rule ccontr)
  assume  $\neg$  ?thesis
  hence  $\text{Re } B = \text{Re } A \vee \text{Re } B = - \text{Re } A$ 
  using power2-eq-iff by blast
  hence  $A * (\text{Re } x)^2 + A * 2 * \text{Re } x + A = 0 \vee A * (\text{Re } x)^2 - A * 2 * \text{Re } x + A = 0$ 
  using eq
  by auto
  hence  $A * ((\text{Re } x)^2 + 2 * \text{Re } x + 1) = 0 \vee A * ((\text{Re } x)^2 - 2 * \text{Re } x + 1) = 0$ 
  by (simp add: field-simps)
  hence  $(\text{Re } x)^2 + 2 * \text{Re } x + 1 = 0 \vee (\text{Re } x)^2 - 2 * \text{Re } x + 1 = 0$ 
  using  $\langle A \neq 0 \rangle$ 
  by simp
  hence  $(\text{Re } x + 1)^2 = 0 \vee (\text{Re } x - 1)^2 = 0$ 
  by (simp add: power2-sum power2-diff field-simps)
  hence  $\text{Re } x = -1 \vee \text{Re } x = 1$ 
  by auto
  thus False
  using  $\langle \text{is-real } x \rangle$  of-complex x  $\in$  unit-disc
  by (auto simp add: cmod-eq-Re)
qed
ultimately
show intersects-x-axis H
  using intersects-x-axis-mk-circline
  using *
  by auto
next
assume intersects-x-axis H
  hence  $(\text{Re } B)^2 > (\text{Re } A)^2$ 
  using * False
  using intersects-x-axis-mk-circline
  by simp
  hence discr:  $(2 * \text{Re } B)^2 - 4 * A * A > 0$ 
  by (simp add: power2-eq-square)
  then obtain x1 x2 where
  eqs:  $A * x1^2 + 2 * \text{Re } B * x1 + A = 0 \wedge A * x2^2 + 2 * \text{Re } B * x2 + A = 0 \wedge x1 \neq x2$ 
  using discriminant-pos-ex[OF  $\langle A \neq 0 \rangle$ , of 2 * Re B A]

```

```

    using discrim-def[of A 2 * Re B A]
  by auto
hence x1 * x2 = 1
  using viette2[OF ‹A ≠ 0›, of 2 * Re B A x1 x2] discr ‹A ≠ 0›
  by auto
have abs x1 ≠ 1 abs x2 ≠ 1
  using eqs discr ‹x1 * x2 = 1›
  by (auto simp add: abs-if power2-eq-square)
hence abs x1 < 1 ∨ abs x2 < 1
  using ‹x1 * x2 = 1›
  by (smt mult-le-cancel-left1 mult-minus-right)
thus ∃ x ∈ unit-disc. x ∈ circline-set H ∩ circline-set x-axis
  using x-axis-intersection-equation[OF *(1-2), of x1]
  using x-axis-intersection-equation[OF *(1-2), of x2]
  using eqs
  by auto
qed
qed
qed

```

8.3 Check if a Poincaré line intersects the y-axis

definition *intersects-y-axis-cmat* :: *complex-mat* ⇒ *bool* **where**
 [simp]: *intersects-y-axis-cmat* H = (let (A, B, C, D) = H in A = 0 ∨ (Im B)² > (Re A)²)

lift-definition *intersects-y-axis-clmat* :: *circline-mat* ⇒ *bool* **is** *intersects-y-axis-cmat*
 done

lift-definition *intersects-y-axis* :: *circline* ⇒ *bool* **is** *intersects-y-axis-clmat*

proof (transfer)

```

fix H1 H2
assume hh: hermitean H1 ∧ H1 ≠ mat-zero and hermitean H2 ∧ H2 ≠ mat-zero
obtain A1 B1 C1 D1 A2 B2 C2 D2 where *: H1 = (A1, B1, C1, D1) H2 = (A2, B2, C2, D2)
  by (cases H1, cases H2, auto)
assume circline-eq-cmat H1 H2
then obtain k where k: k ≠ 0 ∧ H2 = cor k *sm H1
  by auto
show intersects-y-axis-cmat H1 = intersects-y-axis-cmat H2
proof -
  have k ≠ 0 ⇒ (Re A1)2 < (Im B1)2 ↔ (k * Re A1)2 < (k * Im B1)2
    by (smt mult-strict-left-mono power2-eq-square semiring-normalization-rules(13) zero-less-power2)
  thus ?thesis
    using * k
    by auto
qed
qed

```

lemma *intersects-x-axis-intersects-y-axis* [simp]:

shows *intersects-x-axis* (moebius-circline (moebius-rotation (pi/2)) H) ↔ *intersects-y-axis* H
unfolding moebius-rotation-def moebius-similarity-def
by simp (transfer, transfer, auto simp add: mat-adj-def mat-cnj-def)

lemma *intersects-y-axis-iff*:

assumes *is-poincare-line* H
shows (∃ y ∈ unit-disc. y ∈ circline-set H ∩ circline-set y-axis) ↔ *intersects-y-axis* H (**is** ?lhs ↔ ?rhs)

proof -

```

let ?R = moebius-rotation (pi / 2)
let ?H' = moebius-circline ?R H
have 1: is-poincare-line ?H'
  using assms
  using unit-circle-fix-preserve-is-poincare-line[OF - assms, of ?R]
  by simp

```

show ?thesis

proof

assume ?lhs

```

then obtain y where y ∈ unit-disc y ∈ circline-set H ∩ circline-set y-axis
  by auto
hence moebius-pt ?R y ∈ unit-disc ∧ moebius-pt ?R y ∈ circline-set ?H' ∩ circline-set x-axis
  using rotation-pi-2-y-axis
by (metis Int-iff circline-set-moebius-circline-E moebius-circline-comp-inv-left moebius-pt-comp-inv-left unit-disc-fix-discI
unit-disc-fix-rotation)
thus ?rhs
  using intersects-x-axis-iff[OF 1]
  using intersects-x-axis-intersects-y-axis[of H]
  by auto
next
assume intersects-y-axis H
hence intersects-x-axis ?H'
  using intersects-x-axis-intersects-y-axis[of H]
  by simp
then obtain x where *: x ∈ unit-disc x ∈ circline-set ?H' ∩ circline-set x-axis
  using intersects-x-axis-iff[OF 1]
  by auto
let ?y = moebius-pt (-?R) x
have ?y ∈ unit-disc ∧ ?y ∈ circline-set H ∩ circline-set y-axis
  using * rotation-pi-2-y-axis[symmetric]
  by (metis Int-iff circline-set-moebius-circline-E moebius-pt-comp-inv-left moebius-rotation-uminus uminus-moebius-def
unit-disc-fix-discI unit-disc-fix-rotation)
thus ?lhs
  by auto
qed
qed

```

8.4 Intersection point of a Poincaré line with the x-axis in the unit disc

definition *calc-x-axis-intersection-cvec* :: *complex* ⇒ *complex* ⇒ *complex-vec* **where**

```

[simp]: calc-x-axis-intersection-cvec A B =
  (let discr = (Re B)2 - (Re A)2 in
    (-Re(B) + sgn (Re B) * sqrt(discr), A))

```

definition *calc-x-axis-intersection-cmat-cvec* :: *complex-mat* ⇒ *complex-vec* **where** [simp]:

```

calc-x-axis-intersection-cmat-cvec H =
  (let (A, B, C, D) = H in
    if A ≠ 0 then
      calc-x-axis-intersection-cvec A B
    else
      (0, 1)
  )

```

lift-definition *calc-x-axis-intersection-clmat-hcoords* :: *circline-mat* ⇒ *complex-homo-coords* **is** *calc-x-axis-intersection-cmat-cvec*
by (auto split: if-split-asm)

lift-definition *calc-x-axis-intersection* :: *circline* ⇒ *complex-homo* **is** *calc-x-axis-intersection-clmat-hcoords*

proof transfer

fix H1 H2

assume *: hermitean H1 ∧ H1 ≠ mat-zero hermitean H2 ∧ H2 ≠ mat-zero

obtain A1 B1 C1 D1 A2 B2 C2 D2 **where** hh: H1 = (A1, B1, C1, D1) H2 = (A2, B2, C2, D2)

by (cases H1, cases H2, auto)

assume circline-eq-cmat H1 H2

then obtain k **where** k: k ≠ 0 H2 = cor k *_{sm} H1

by auto

have *calc-x-axis-intersection-cvec* A1 B1 ≈_v *calc-x-axis-intersection-cvec* A2 B2

using hh k

apply simp

apply (rule-tac x=cor k in exI)

apply auto

apply (simp add: sgn-mult power-mult-distrib)

apply (subst right-diff-distrib[symmetric])

apply (subst real-sqrt-mult)

by (simp add: real-sgn-eq right-diff-distrib)

thus calc-x-axis-intersection-cmat-cvec H1 \approx_v
 calc-x-axis-intersection-cmat-cvec H2
 using hh k
 by (auto simp del: calc-x-axis-intersection-cvec-def)

qed

lemma calc-x-axis-intersection-in-unit-disc:
 assumes is-poincare-line H intersects-x-axis H
 shows calc-x-axis-intersection H \in unit-disc

proof (cases is-line H)
 case True
 thus ?thesis
 using assms
 unfolding unit-disc-def disc-def
 by simp (transfer, transfer, auto simp add: vec-cnj-def)

next

case False

thus ?thesis
 using assms

unfolding unit-disc-def disc-def

proof (simp, transfer, transfer)

fix H

assume hh: hermitean H \wedge H \neq mat-zero

then obtain A B D where *: H = (A, B, cnj B, D) is-real A is-real D

using hermitean-elems

by (cases H) blast

assume is-poincare-line-cmat H

hence *: H = (A, B, cnj B, A) is-real A

using *

by auto

assume \neg circline-A0-cmat H

hence A \neq 0

using *

by simp

assume intersects-x-axis-cmat H

hence $(\operatorname{Re} B)^2 > (\operatorname{Re} A)^2$

using * $\langle A \neq 0 \rangle$

by (auto simp add: power2-eq-square complex.expand)

hence $\operatorname{Re} B \neq 0$

by auto

have $\operatorname{Re} A \neq 0$

using $\langle \text{is-real } A \rangle \langle A \neq 0 \rangle$

by (auto simp add: complex.expand)

have $\sqrt{(\operatorname{Re} B)^2 - (\operatorname{Re} A)^2} < \sqrt{(\operatorname{Re} B)^2}$

using $\langle \operatorname{Re} A \neq 0 \rangle$

by (subst real-sqrt-less-iff) auto

also have ... = $\operatorname{sgn} (\operatorname{Re} B) * (\operatorname{Re} B)$

by (smt mult-minus-right nonzero-eq-divide-eq real-sgn-eq real-sqrt-abs)

finally

have 1: $\sqrt{(\operatorname{Re} B)^2 - (\operatorname{Re} A)^2} < \operatorname{sgn} (\operatorname{Re} B) * (\operatorname{Re} B)$

.

have 2: $(\operatorname{Re} B)^2 - (\operatorname{Re} A)^2 < \operatorname{sgn} (\operatorname{Re} B) * (\operatorname{Re} B) * \sqrt{(\operatorname{Re} B)^2 - (\operatorname{Re} A)^2}$

using $\langle (\operatorname{Re} B)^2 > (\operatorname{Re} A)^2 \rangle$

using mult-strict-right-mono[OF 1, of $\sqrt{(\operatorname{Re} B)^2 - (\operatorname{Re} A)^2}$]

by simp

have 3: $(\operatorname{Re} B)^2 - 2 * \operatorname{sgn} (\operatorname{Re} B) * \operatorname{Re} B * \sqrt{(\operatorname{Re} B)^2 - (\operatorname{Re} A)^2} + (\operatorname{Re} B)^2 - (\operatorname{Re} A)^2 < (\operatorname{Re} A)^2$

```

using mult-strict-left-mono[OF 2, of 2]
by (simp add: field-simps)

have (sgn (Re B))2 = 1
  using ⟨Re B ≠ 0⟩
  by (simp add: sgn-if)

hence (-Re B + sgn (Re B) * sqrt((Re B)2 - (Re A)2))2 < (Re A)2
  using ⟨(Re B)2 > (Re A)2⟩ 3
  by (simp add: power2-diff field-simps)

thus in-ocircline-cmat-cvec unit-circle-cmat (calc-x-axis-intersection-cmat-cvec H)
  using * ⟨(Re B)2 > (Re A)2⟩
  by (auto simp add: vec-cnj-def power2-eq-square split: if-split-asm)
qed
qed

lemma calc-x-axis-intersection:
  assumes is-poincare-line H and intersects-x-axis H
  shows calc-x-axis-intersection H ∈ circline-set H ∩ circline-set x-axis
proof (cases is-line H)
case True
  thus ?thesis
    using assms
    unfolding circline-set-def
    by simp (transfer, transfer, auto simp add: vec-cnj-def)
next
case False
  thus ?thesis
    using assms
    unfolding circline-set-def
proof (simp, transfer, transfer)
  fix H
  assume hh: hermitean H ∧ H ≠ mat-zero
  then obtain A B D where *: H = (A, B, cnj B, D) is-real A is-real D
    using hermitean-elems
    by (cases H) blast
  assume is-poincare-line-cmat H
  hence *: H = (A, B, cnj B, A) is-real A
    using *
    by auto
  assume ¬ circline-A0-cmat H
  hence A ≠ 0
    using *
    by auto

  assume intersects-x-axis-cmat H
  hence (Re B)2 > (Re A)2
    using * ⟨A ≠ 0⟩
    by (auto simp add: power2-eq-square complex.expand)

  hence Re B ≠ 0
    by auto

  show on-circline-cmat-cvec H (calc-x-axis-intersection-cmat-cvec H) ∧
    on-circline-cmat-cvec x-axis-cmat (calc-x-axis-intersection-cmat-cvec H) (is ?P1 ∧ ?P2)
proof
  show on-circline-cmat-cvec H (calc-x-axis-intersection-cmat-cvec H)
  proof (cases circline-A0-cmat H)
  case True
    thus ?thesis
      using * ⟨is-poincare-line-cmat H⟩ ⟨intersects-x-axis-cmat H⟩
      by (simp add: vec-cnj-def)
  next
  case False

```

let $?x = \text{calc-x-axis-intersection-cvec } A \ B$
let $?nom = \text{fst } ?x$ **and** $?den = \text{snd } ?x$
have $x: ?x = (?nom, ?den)$
by *simp*

hence *on-circline-cmat-cvec H (calc-x-axis-intersection-cvec A B)*
proof (*subst **, *subst x*, *subst on-circline-cmat-cvec-circline-equation*)
have $(\text{sgn}(\text{Re } B))^2 = 1$
using $\langle \text{Re } B \neq 0 \rangle$ *sgn-pos zero-less-power2* **by** *fastforce*
have $(\text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 = (\text{Re } B)^2 - (\text{Re } A)^2$
using $\langle (\text{Re } B)^2 > (\text{Re } A)^2 \rangle$,
by *simp*

have $(-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 =$
 $(-\text{Re } B)^2 + (\text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)$
by (*simp add: power2-diff*)
also have $\dots = (\text{Re } B) * (\text{Re } B) + (\text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)$
by (*simp add: power2-eq-square*)
also have $\dots = (\text{Re } B) * (\text{Re } B) + (\text{sgn}(\text{Re } B))^2 * (\text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)$
by (*simp add: power-mult-distrib*)
also have $\dots = (\text{Re } B) * (\text{Re } B) + (\text{Re } B)^2 - (\text{Re } A)^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)$
using $\langle (\text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 = (\text{Re } B)^2 - (\text{Re } A)^2 \rangle$, $\langle (\text{sgn}(\text{Re } B))^2 = 1 \rangle$
by *simp*
finally have $(-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 =$
 $(\text{Re } B) * (\text{Re } B) + (\text{Re } B)^2 - (\text{Re } A)^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)$
by *simp*

have *is-real ?nom is-real ?den*
using $\langle \text{is-real } A \rangle$
by *simp+*
hence $\text{cnj } (?nom) = ?nom$ $\text{cnj } (?den) = ?den$
by (*simp add: eq-cnj-iff-real*)
hence $A * ?nom * (\text{cnj } (?nom)) + B * ?den * (\text{cnj } (?nom)) + (\text{cnj } B) * (\text{cnj } (?den)) * ?nom + A * ?den * (\text{cnj } (?den))$
 $= A * ?nom * ?nom + B * ?den * ?nom + (\text{cnj } B) * ?den * ?nom + A * ?den * ?den$
by *auto*
also have $\dots = A * ?nom * ?nom + (B + (\text{cnj } B)) * ?den * ?nom + A * ?den * ?den$
by (*simp add: field-simps*)
also have $\dots = A * ?nom * ?nom + 2 * (\text{Re } B) * ?den * ?nom + A * ?den * ?den$
by (*simp add: complex-add-cnj*)
also have $\dots = A * ?nom^2 + 2 * (\text{Re } B) * ?den * ?nom + A * ?den * ?den$
by (*simp add: power2-eq-square*)
also have $\dots = A * (-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2$
 $+ 2 * (\text{Re } B) * A * (-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)) + A * A * A$
unfolding *calc-x-axis-intersection-cvec-def*
by *auto*
also have $\dots = A * ((\text{Re } B) * (\text{Re } B) + (\text{Re } B)^2 - (\text{Re } A)^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))$
 $+ 2 * (\text{Re } B) * A * (-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)) + A * A * A$
using $\langle (-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))^2 =$
 $(\text{Re } B) * (\text{Re } B) + (\text{Re } B)^2 - (\text{Re } A)^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2) \rangle$,
by *simp*
also have $\dots = A * ((\text{Re } B) * (\text{Re } B) + (\text{Re } B)^2 - A^2 - 2 * (\text{Re } B) * \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2))$
 $+ 2 * (\text{Re } B) * A * (-(\text{Re } B) + \text{sgn}(\text{Re } B) * \text{sqrt}((\text{Re } B)^2 - (\text{Re } A)^2)) + A * A * A$
using $\langle \text{is-real } A \rangle$
by *simp*
also have $\dots = 0$
apply (*simp add: field-simps*)
by (*simp add: power2-eq-square*)
finally have $A * ?nom * (\text{cnj } (?nom)) + B * ?den * (\text{cnj } (?nom)) + (\text{cnj } B) * (\text{cnj } (?den)) * ?nom + A * ?den * (\text{cnj } (?den)) = 0$
by *simp*
thus *circline-equation A B (cnj B) A ?nom ?den*
by *simp*

qed
thus *?thesis*


```

    using * ⟨is-poincare-line-cmat H⟩ ⟨intersects-x-axis-cmat H⟩
    by (simp add: vec-cnj-def)
  qed
next
show on-circline-cmat-cvec x-axis-cmat (calc-x-axis-intersection-cmat-cvec H)
  using * ⟨is-poincare-line-cmat H⟩ ⟨intersects-x-axis-cmat H⟩ ⟨is-real A⟩
  using eq-cnj-iff-real[of A]
  by (simp add: vec-cnj-def)
qed
qed
qed

```

```

lemma unique-calc-x-axis-intersection:
  assumes is-poincare-line H and H ≠ x-axis
  assumes x ∈ unit-disc and x ∈ circline-set H ∩ circline-set x-axis
  shows x = calc-x-axis-intersection H
proof -
  have *: intersects-x-axis H
  using assms
  using intersects-x-axis-iff[OF assms(1)]
  by auto
show x = calc-x-axis-intersection H
  using calc-x-axis-intersection[OF assms(1) *]
  using calc-x-axis-intersection-in-unit-disc[OF assms(1) *]
  using assms
  using unique-is-poincare-line[of x calc-x-axis-intersection H H x-axis]
  by auto
qed

```

8.5 Check if an h-line intersects the positive part of the x-axis

definition *intersects-x-axis-positive-cmat* :: complex-mat \Rightarrow bool **where**
 [simp]: *intersects-x-axis-positive-cmat* H = (let (A, B, C, D) = H in Re A \neq 0 \wedge Re B / Re A < -1)

lift-definition *intersects-x-axis-positive-clmat* :: circline-mat \Rightarrow bool **is** *intersects-x-axis-positive-cmat*
 done

lift-definition *intersects-x-axis-positive* :: circline \Rightarrow bool **is** *intersects-x-axis-positive-clmat*

```

proof (transfer)
  fix H1 H2
  assume hh: hermitean H1  $\wedge$  H1  $\neq$  mat-zero and hermitean H2  $\wedge$  H2  $\neq$  mat-zero
  obtain A1 B1 C1 D1 A2 B2 C2 D2 where *: H1 = (A1, B1, C1, D1) H2 = (A2, B2, C2, D2)
  by (cases H1, cases H2, auto)
  assume circline-eq-cmat H1 H2
  then obtain k where k  $\neq$  0  $\wedge$  H2 = cor k *sm H1
  by auto
  thus intersects-x-axis-positive-cmat H1 = intersects-x-axis-positive-cmat H2
  using *
  by simp
qed

```

```

lemma intersects-x-axis-positive-mk-circline:
  assumes is-real A and A  $\neq$  0  $\vee$  B  $\neq$  0
  shows intersects-x-axis-positive (mk-circline A B (cnj B) A)  $\longleftrightarrow$  Re B / Re A < -1
proof -

```

```

let ?H = (A, B, (cnj B), A)
have hermitean ?H
  using ⟨is-real A⟩
  unfolding hermitean-def mat-adj-def mat-cnj-def
  using eq-cnj-iff-real
  by auto
moreover
have ?H  $\neq$  mat-zero
  using assms
  by auto
ultimately

```

show ?thesis
 by (transfer, transfer, auto simp add: Let-def)
 qed

lemma intersects-x-axis-positive-intersects-x-axis [simp]:

assumes intersects-x-axis-positive H
 shows intersects-x-axis H

proof-

have $\bigwedge a aa. \llbracket \text{Re } a \neq 0; \text{Re } aa / \text{Re } a < -1; \neg (\text{Re } a)^2 < (\text{Re } aa)^2 \rrbracket \implies aa = 0 \wedge a = 0$
 by (smt less-divide-eq-1-pos one-less-power pos2 power2-minus power-divide zero-less-power2)
 thus ?thesis
 using assms
 apply transfer
 apply transfer
 apply (auto simp add: hermitean-def mat-adj-def mat-cnj-def)
 done

qed

lemma add-less-abs-positive-iff:

fixes a b :: real
 assumes abs b < abs a
 shows $a + b > 0 \iff a > 0$
 using assms
 by auto

lemma calc-x-axis-intersection-positive-abs':

fixes A B :: real
 assumes $B^2 > A^2$ and $A \neq 0$
 shows $\text{abs} (\text{sgn}(B) * \text{sqrt}(B^2 - A^2) / A) < \text{abs}(-B/A)$

proof-

from assms have $B \neq 0$
 by auto

have $B^2 - A^2 < B^2$
 using $\langle A \neq 0 \rangle$
 by auto

hence $\text{sqrt}(B^2 - A^2) < \text{abs } B$
 using real-sqrt-less-iff[of $B^2 - A^2$ B^2]
 by simp

thus ?thesis
 using assms $\langle B \neq 0 \rangle$
 by (simp add: abs-mult divide-strict-right-mono)

qed

lemma calc-intersect-x-axis-positive-lemma:

assumes $B^2 > A^2$ and $A \neq 0$
 shows $(-B + \text{sgn } B * \text{sqrt}(B^2 - A^2)) / A > 0 \iff -B/A > 1$

proof-

have $(-B + \text{sgn } B * \text{sqrt}(B^2 - A^2)) / A = -B / A + (\text{sgn } B * \text{sqrt}(B^2 - A^2)) / A$
 using assms
 by (simp add: field-simps)

moreover

have $-B / A + (\text{sgn } B * \text{sqrt}(B^2 - A^2)) / A > 0 \iff -B / A > 0$
 using add-less-abs-positive-iff[OF calc-x-axis-intersection-positive-abs'[OF assms]]
 by simp

moreover

hence $(B/A)^2 > 1$
 using assms

by (simp add: power-divide)

hence $B/A > 1 \vee B/A < -1$

by (smt one-power2 pos2 power2-minus power-0 power-strict-decreasing zero-power2)

hence $-B / A > 0 \iff -B / A > 1$

by auto

ultimately

show ?thesis

```

using assms
by auto
qed

lemma intersects-x-axis-positive-iff':
  assumes is-poincare-line H
  shows intersects-x-axis-positive H  $\longleftrightarrow$ 
    calc-x-axis-intersection H  $\in$  unit-disc  $\wedge$  calc-x-axis-intersection H  $\in$  circline-set H  $\cap$  positive-x-axis (is ?lhs  $\longleftrightarrow$ 
?rhs)
proof
  let ?x = calc-x-axis-intersection H
  assume ?lhs
  hence ?x  $\in$  circline-set x-axis ?x  $\in$  circline-set H ?x  $\in$  unit-disc
    using calc-x-axis-intersection-in-unit-disc[OF assms] calc-x-axis-intersection[OF assms]
    by auto
  moreover
  have Re (to-complex ?x) > 0
    using <?lhs> assms
  proof (transfer, transfer)
    fix H
    assume hh: hermitean H  $\wedge$  H  $\neq$  mat-zero
    obtain A B C D where *: H = (A, B, C, D)
      by (cases H, auto)
    assume intersects-x-axis-positive-cmat H
    hence **: Re B / Re A < - 1 Re A  $\neq$  0
      using *
      by auto
    have (Re B)2 > (Re A)2
      using **
      by (smt divide-less-eq-1-neg divide-minus-left less-divide-eq-1-pos real-sqrt-abs real-sqrt-less-iff right-inverse-eq)
    have is-real A A  $\neq$  0
      using hh hermitean-elems * <Re A  $\neq$  0> complex.expand[of A 0]
      by auto
    have (cmod B)2 > (cmod A)2
      using <(Re B)2 > (Re A)22 - (Re A)2)) / Re A
      using calc-intersect-x-axis-positive-lemma[of Re A Re B] ** <(Re B)2 > (Re A)22 - (Re A)2)) - cor (Re B)) / A) = (sgn (Re B) * sqrt ((Re B)2
- (Re A)2) - Re B) / Re A
      using <is-real A> <A = D>
      by (metis (no-types, lifting) Re-complex-of-real complex-of-real-Re of-real-diff of-real-divide of-real-mult)
    thus 0 < Re (to-complex-cvec (calc-x-axis-intersection-cmat-cvec H))
      using * hh ** *** <(cmod B)2 > (cmod A)22 > (Re A)2\neq 0> <A = D>
      by simp
  qed
  ultimately
  show ?rhs
    unfolding positive-x-axis-def
    by auto
next
  let ?x = calc-x-axis-intersection H
  assume ?rhs
  hence Re (to-complex ?x) > 0 ?x  $\neq$   $\infty_h$  ?x  $\in$  circline-set x-axis ?x  $\in$  unit-disc ?x  $\in$  circline-set H
    unfolding positive-x-axis-def
    by auto
  hence intersects-x-axis H
    using intersects-x-axis-iff[OF assms]
    by auto

```

```

thus ?lhs
  using ⟨Re (to-complex ?x) > 0⟩ assms
proof (transfer, transfer)
  fix H
  assume hh: hermitean H ∧ H ≠ mat-zero
  obtain A B C D where *: H = (A, B, C, D)
    by (cases H, auto)
  assume 0 < Re (to-complex-cvec (calc-x-axis-intersection-cmat-cvec H)) intersects-x-axis-cmat H is-poincare-line-cmat
H
  hence **: A ≠ 0 0 < Re ((cor (sgn (Re B)) * cor (sqrt ((Re B)2 - (Re A)2)) - cor (Re B)) / A) A = D is-real
A (Re B)2 > (Re A)2
  using * hh hermitean-elems
  by (auto split: if-split-asm)

  have Re A ≠ 0
  using complex.expand[of A 0] ⟨A ≠ 0⟩ ⟨is-real A⟩
  by auto

  have Re ((cor (sgn (Re B)) * cor (sqrt ((Re B)2 - (Re D)2)) - cor (Re B)) / D) = (sgn (Re B) * sqrt ((Re B)2
- (Re D)2) - Re B) / Re D
  using ⟨is-real A⟩ ⟨A = D⟩
  by (metis (no-types, lifting) Re-complex-of-real complex-of-real-Re of-real-diff of-real-divide of-real-mult)

  thus intersects-x-axis-positive-cmat H
  using * ** ⟨Re A ≠ 0⟩
  using calc-intersect-x-axis-positive-lemma[of Re A Re B]
  by simp
qed
qed

```

```

lemma intersects-x-axis-positive-iff:
assumes is-poincare-line H and H ≠ x-axis
shows intersects-x-axis-positive H ⟷
  (∃ x. x ∈ unit-disc ∧ x ∈ circline-set H ∩ positive-x-axis) (is ?lhs ⟷ ?rhs)
proof
  assume ?lhs
  thus ?rhs
  using intersects-x-axis-positive-iff'[OF assms(1)]
  by auto
next
  assume ?rhs
  then obtain x where x ∈ unit-disc x ∈ circline-set H ∩ positive-x-axis
  by auto
  thus ?lhs
  using unique-calc-x-axis-intersection[OF assms, of x]
  using intersects-x-axis-positive-iff'[OF assms(1)]
  unfolding positive-x-axis-def
  by auto
qed

```

8.6 Check if an h-line intersects the positive part of the y-axis

definition intersects-y-axis-positive-cmat :: complex-mat ⇒ bool **where**
 [*simp*]: intersects-y-axis-positive-cmat H = (let (A, B, C, D) = H in Re A ≠ 0 ∧ Im B / Re A < -1)

lift-definition intersects-y-axis-positive-clmat :: circline-mat ⇒ bool **is** intersects-y-axis-positive-cmat
 done

lift-definition intersects-y-axis-positive :: circline ⇒ bool **is** intersects-y-axis-positive-clmat

```

proof (transfer)
  fix H1 H2
  assume hh: hermitean H1 ∧ H1 ≠ mat-zero and hermitean H2 ∧ H2 ≠ mat-zero
  obtain A1 B1 C1 D1 A2 B2 C2 D2 where *: H1 = (A1, B1, C1, D1) H2 = (A2, B2, C2, D2)
  by (cases H1, cases H2, auto)
  assume circline-eq-cmat H1 H2
  then obtain k where k ≠ 0 ∧ H2 = cor k *sm H1

```

by *auto*
 thus *intersects-y-axis-positive-cmat* $H1 = intersects-y-axis-positive-cmat$ $H2$
 using ***
 by *simp*
 qed

lemma *intersects-x-axis-positive-intersects-y-axis-positive* [*simp*]:
shows *intersects-x-axis-positive* (*moebius-circline* (*moebius-rotation* $(-\pi/2)$) H) $\longleftrightarrow intersects-y-axis-positive$ H
 using *hermitean-elems*
unfolding *moebius-rotation-def* *moebius-similarity-def*
by *simp* (*transfer*, *transfer*, *auto* *simp* *add: mat-adj-def* *mat-cnj-def*)

lemma *intersects-y-axis-positive-iff*:
assumes *is-poincare-line* H $H \neq y\text{-axis}$
shows $(\exists y \in \text{unit-disc. } y \in \text{circline-set } H \cap \text{positive-y-axis}) \longleftrightarrow intersects-y-axis-positive$ H (**is** *?lhs* \longleftrightarrow *?rhs*)

proof–
let $?R = \text{moebius-rotation } (-\pi / 2)$
let $?H' = \text{moebius-circline } ?R$ H
have *1: is-poincare-line* $?H'$
 using *assms*
 using *unit-circle-fix-preserve-is-poincare-line*[*OF - assms(1)*, *of ?R*]
by *simp*

have *2: moebius-circline* $?R$ $H \neq x\text{-axis}$

proof (*rule ccontr*)
assume $\neg ?thesis$
hence $H = \text{moebius-circline } (\text{moebius-rotation } (\pi/2))$ $x\text{-axis}$
 using *moebius-circline-comp-inv-left*[*of ?R* H]
by *auto*
thus *False*
 using $\langle H \neq y\text{-axis} \rangle$
by *auto*

qed

show *?thesis*

proof
assume *?lhs*
then obtain y **where** $y \in \text{unit-disc } y \in \text{circline-set } H \cap \text{positive-y-axis}$
by *auto*
hence *moebius-pt* $?R$ $y \in \text{unit-disc } \text{moebius-pt } ?R$ $y \in \text{circline-set } ?H' \cap \text{positive-x-axis}$
 using *rotation-minus-pi-2-positive-y-axis*
by *auto*
thus *?rhs*
 using *intersects-x-axis-positive-iff*[*OF 1 2*]
 using *intersects-x-axis-positive-intersects-y-axis-positive*[*of H*]
by *auto*

next

assume *intersects-y-axis-positive* H
hence *intersects-x-axis-positive* $?H'$
 using *intersects-x-axis-positive-intersects-y-axis-positive*[*of H*]
by *simp*

then obtain x **where** $*$: $x \in \text{unit-disc } x \in \text{circline-set } ?H' \cap \text{positive-x-axis}$
 using *intersects-x-axis-positive-iff*[*OF 1 2*]
by *auto*

let $?y = \text{moebius-pt } (-?R)$ x

have $?y \in \text{unit-disc} \wedge ?y \in \text{circline-set } H \cap \text{positive-y-axis}$
 using *** *rotation-minus-pi-2-positive-y-axis*[*symmetric*]

by (*metis* *Int-iff* *circline-set-moebius-circline-E* *image-eqI* *moebius-pt-comp-inv-image-left* *moebius-rotation-uminus*

uminus-moebius-def *unit-disc-fix-discI* *unit-disc-fix-rotation*)

thus *?lhs*

by *auto*

qed

qed

8.7 Position of the intersection point in the unit disc

Check if the intersection point of one h-line with the x-axis is located more outward the edge of the disc than the intersection point of another h-line.

definition *outward-cmat* :: *complex-mat* \Rightarrow *complex-mat* \Rightarrow *bool* **where**
 [simp]: *outward-cmat* *H1* *H2* = (let (*A1*, *B1*, *C1*, *D1*) = *H1*; (*A2*, *B2*, *C2*, *D2*) = *H2*
 in $-\text{Re } B1 / \text{Re } A1 \leq -\text{Re } B2 / \text{Re } A2$)

lift-definition *outward-clmat* :: *circline-mat* \Rightarrow *circline-mat* \Rightarrow *bool* **is** *outward-cmat*
done

lift-definition *outward* :: *circline* \Rightarrow *circline* \Rightarrow *bool* **is** *outward-clmat*

apply *transfer*

apply *simp*

apply (*case-tac circline-mat1*, *case-tac circline-mat2*, *case-tac circline-mat3*, *case-tac circline-mat4*)

apply *simp*

apply (*erule-tac exE*)⁺

apply (*erule-tac conjE*)⁺

apply *simp*

done

lemma *outward-mk-circline*:

assumes *is-real A1* **and** *is-real A2* **and** $A1 \neq 0 \vee B1 \neq 0$ **and** $A2 \neq 0 \vee B2 \neq 0$

shows *outward* (*mk-circline A1 B1 (cnj B1) A1*) (*mk-circline A2 B2 (cnj B2) A2*) $\longleftrightarrow -\text{Re } B1 / \text{Re } A1 \leq -\text{Re } B2 / \text{Re } A2$

proof–

let $?H1 = (A1, B1, (\text{cnj } B1), A1)$

let $?H2 = (A2, B2, (\text{cnj } B2), A2)$

have *hermitean ?H1* *hermitean ?H2*

using $\langle \text{is-real } A1 \rangle \langle \text{is-real } A2 \rangle$

unfolding *hermitean-def mat-adj-def mat-cnj-def*

using *eq-cnj-iff-real*

by *auto*

moreover

have $?H1 \neq \text{mat-zero}$ $?H2 \neq \text{mat-zero}$

using *assms*

by *auto*

ultimately

show *?thesis*

by (*transfer*, *transfer*, *auto simp add: Let-def*)

qed

lemma *calc-x-axis-intersection-fun-mono*:

fixes *x1 x2* :: *real*

assumes $x1 > 1$ **and** $x2 > x1$

shows $x1 - \text{sqrt}(x1^2 - 1) > x2 - \text{sqrt}(x2^2 - 1)$

using *assms*

proof–

have *: $\text{sqrt}(x1^2 - 1) + \text{sqrt}(x2^2 - 1) > 0$

using *assms*

by (*smt one-less-power pos2 real-sqrt-gt-zero*)

have $\text{sqrt}(x1^2 - 1) < x1$

using *real-sqrt-less-iff*[*of* $x1^2 - 1$ $x1^2$] $\langle x1 > 1 \rangle$

by *auto*

moreover

have $\text{sqrt}(x2^2 - 1) < x2$

using *real-sqrt-less-iff*[*of* $x2^2 - 1$ $x2^2$] $\langle x1 > 1 \rangle \langle x2 > x1 \rangle$

by *auto*

ultimately

have $\text{sqrt}(x1^2 - 1) + \text{sqrt}(x2^2 - 1) < x1 + x2$

by *simp*

hence $(x1 + x2) / (\text{sqrt}(x1^2 - 1) + \text{sqrt}(x2^2 - 1)) > 1$

using *

using *less-divide-eq-1-pos*[*of* $\text{sqrt}(x1^2 - 1) + \text{sqrt}(x2^2 - 1)$ $x1 + x2$]

by *simp*

hence $(x2^2 - x1^2) / (\text{sqrt}(x1^2 - 1) + \text{sqrt}(x2^2 - 1)) > x2 - x1$

using $\langle x2 > x1 \rangle$

```

using mult-less-cancel-left-pos[of x2 - x1 1 (x2 + x1) / (sqrt(x12 - 1) + sqrt(x22 - 1))]
by (simp add: power2-eq-square field-simps)
moreover
have (x22 - x12) = (sqrt(x12 - 1) + sqrt(x22 - 1)) * ((sqrt(x22 - 1) - sqrt(x12 - 1)))
using ⟨x1 > 1⟩ ⟨x2 > x1⟩
by (simp add: field-simps)
ultimately
have sqrt(x22 - 1) - sqrt(x12 - 1) > x2 - x1
using *
by simp
thus ?thesis
by simp
qed

```

lemma calc-x-axis-intersection-mono:

```

fixes a1 b1 a2 b2 :: real
assumes -b1/a1 > 1 and a1 ≠ 0 and -b2/a2 ≥ -b1/a1 and a2 ≠ 0
shows (-b1 + sgn b1 * sqrt(b12 - a12)) / a1 ≥ (-b2 + sgn b2 * sqrt(b22 - a22)) / a2 (is ?lhs ≥ ?rhs)

```

proof-

```

have ?lhs = -b1/a1 - sqrt((-b1/a1)2 - 1)

```

```

proof (cases b1 > 0)

```

```

case True

```

```

hence a1 < 0

```

```

using assms

```

```

by (smt divide-neg-pos)

```

```

thus ?thesis

```

```

using ⟨b1 > 0⟩ ⟨a1 < 0⟩

```

```

by (simp add: real-sqrt-divide field-simps)

```

next

```

case False

```

```

hence b1 < 0

```

```

using assms

```

```

by (cases b1 = 0) auto

```

```

hence a1 > 0

```

```

using assms

```

```

by (smt divide-pos-neg)

```

```

thus ?thesis

```

```

using ⟨b1 < 0⟩ ⟨a1 > 0⟩

```

```

by (simp add: real-sqrt-divide field-simps)

```

qed

moreover

```

have ?rhs = -b2/a2 - sqrt((-b2/a2)2 - 1)

```

```

proof (cases b2 > 0)

```

```

case True

```

```

hence a2 < 0

```

```

using assms

```

```

by (smt divide-neg-pos)

```

```

thus ?thesis

```

```

using ⟨b2 > 0⟩ ⟨a2 < 0⟩

```

```

by (simp add: real-sqrt-divide field-simps)

```

next

```

case False

```

```

hence b2 < 0

```

```

using assms

```

```

by (cases b2 = 0) auto

```

```

hence a2 > 0

```

```

using assms

```

```

by (smt divide-pos-neg)

```

```

thus ?thesis

```

```

using ⟨b2 < 0⟩ ⟨a2 > 0⟩

```

```

by (simp add: real-sqrt-divide field-simps)

```

qed

ultimately

```

show ?thesis
  using calc-x-axis-intersection-fun-mono[of  $-b1/a1 -b2/a2$ ]
  using assms
  by (cases  $-b1/a1 = -b2/a2$ , auto)
qed

lemma outward:
  assumes is-poincare-line H1 and is-poincare-line H2
  assumes intersects-x-axis-positive H1 and intersects-x-axis-positive H2
  assumes outward H1 H2
  shows Re (to-complex (calc-x-axis-intersection H1))  $\geq$  Re (to-complex (calc-x-axis-intersection H2))
proof -
  have intersects-x-axis H1 intersects-x-axis H2
  using assms
  by auto
  thus ?thesis
  using assms
proof (transfer, transfer)
  fix H1 H2
  assume hh: hermitean H1  $\wedge$  H1  $\neq$  mat-zero hermitean H2  $\wedge$  H2  $\neq$  mat-zero
  obtain A1 B1 C1 D1 A2 B2 C2 D2 where *: H1 = (A1, B1, C1, D1) H2 = (A2, B2, C2, D2)
  by (cases H1, cases H2, auto)
  have is-real A1 is-real A2
  using hermitean-elems * hh
  by auto
  assume 1: intersects-x-axis-positive-cmat H1 intersects-x-axis-positive-cmat H2
  assume 2: intersects-x-axis-cmat H1 intersects-x-axis-cmat H2
  assume 3: is-poincare-line-cmat H1 is-poincare-line-cmat H2
  assume 4: outward-cmat H1 H2
  have A1  $\neq$  0 A2  $\neq$  0
  using *  $\langle$ is-real A1 $\rangle$   $\langle$ is-real A2 $\rangle$  1 complex.expand[of A1 0] complex.expand[of A2 0]
  by auto
  hence (sgn (Re B2) * sqrt ((Re B2)2 - (Re A2)2) - Re B2) / Re A2
   $\leq$  (sgn (Re B1) * sqrt ((Re B1)2 - (Re A1)2) - Re B1) / Re A1
  using calc-x-axis-intersection-mono[of Re B1 Re A1 Re B2 Re A2]
  using 1 4 *
  by simp
  moreover
  have (sgn (Re B2) * sqrt ((Re B2)2 - (Re A2)2) - Re B2) / Re A2 =
  Re ((cor (sgn (Re B2)) * cor (sqrt ((Re B2)2 - (Re A2)2)) - cor (Re B2)) / A2)
  using  $\langle$ is-real A2 $\rangle$   $\langle$ A2  $\neq$  0 $\rangle$ 
  by (simp add: Re-divide-real)
  moreover
  have (sgn (Re B1) * sqrt ((Re B1)2 - (Re A1)2) - Re B1) / Re A1 =
  Re ((cor (sgn (Re B1)) * cor (sqrt ((Re B1)2 - (Re A1)2)) - cor (Re B1)) / A1)
  using  $\langle$ is-real A1 $\rangle$   $\langle$ A1  $\neq$  0 $\rangle$ 
  by (simp add: Re-divide-real)
  ultimately
  show Re (to-complex-cvec (calc-x-axis-intersection-cmat-cvec H2))
   $\leq$  Re (to-complex-cvec (calc-x-axis-intersection-cmat-cvec H1))
  using 2 3  $\langle$ A1  $\neq$  0 $\rangle$   $\langle$ A2  $\neq$  0 $\rangle$  *  $\langle$ is-real A1 $\rangle$   $\langle$ is-real A2 $\rangle$ 
  by (simp del: is-poincare-line-cmat-def intersects-x-axis-cmat-def)
qed
qed

```

8.8 Ideal points and x-axis intersection

```

lemma ideal-points-intersects-x-axis:
  assumes is-poincare-line H and ideal-points H = {i1, i2} and H  $\neq$  x-axis
  shows intersects-x-axis H  $\iff$  Im (to-complex i1) * Im (to-complex i2) < 0
  using assms
proof -
  have i1  $\neq$  i2
  using assms(1) assms(2) ex-poincare-line-points ideal-points-different(1)
  by blast

```



```

have calc-ideal-points H = {i1, i2}
  using assms
  using ideal-points-unique
  by auto

have  $\forall i1 \in \text{calc-ideal-points } H.$ 
   $\forall i2 \in \text{calc-ideal-points } H.$ 
    is-poincare-line H  $\wedge$  H  $\neq$  x-axis  $\wedge$  i1  $\neq$  i2  $\longrightarrow$  (Im (to-complex i1) * Im (to-complex i2) < 0  $\longleftrightarrow$ 
intersects-x-axis H)
proof (transfer, transfer, (rule ballI)+, rule impI, (erule conjE)+, case-tac H, case-tac i1, case-tac i2)
  fix i11 i12 i21 i22 A B C D H i1 i2
  assume H: H = (A, B, C, D) hermitean H H  $\neq$  mat-zero
  assume line: is-poincare-line-cmat H
  assume i1: i1 = (i11, i12) i1  $\in$  calc-ideal-points-cmat-cvec H
  assume i2: i2 = (i21, i22) i2  $\in$  calc-ideal-points-cmat-cvec H
  assume different:  $\neg$  i1  $\approx_v$  i2
  assume not-x-axis:  $\neg$  circline-eq-cmat H x-axis-cmat

  have is-real A is-real D C = conj B
    using H hermitean-elems
    by auto
  have (cmod A)2 < (cmod B)2 A = D
    using line H
    by auto

  let ?discr = sqrt ((cmod B)2 - (Re D)2)
  let ?den = (cmod B)2
  let ?i1 = B * (- D - i * ?discr)
  let ?i2 = B * (- D + i * ?discr)

  have i11 = ?i1  $\vee$  i11 = ?i2 i12 = ?den
    i21 = ?i1  $\vee$  i21 = ?i2 i22 = ?den
    using i1 i2 H line
    by (auto split: if-split-asm)
  hence i: i11 = ?i1  $\wedge$  i21 = ?i2  $\vee$  i11 = ?i2  $\wedge$  i21 = ?i1
    using  $\langle \neg$  i1  $\approx_v$  i2  $\rangle$  i1 i2
    by auto

  have Im (i11 / i12) * Im (i21 / i22) = Im (?i1 / ?den) * Im (?i2 / ?den)
    using i  $\langle$  i12 = ?den  $\rangle$   $\langle$  i22 = ?den  $\rangle$ 
    by auto
  also have ... = Im (?i1) * Im (?i2) / ?den2
    by simp
  also have ... = (Im B * (Im B * (Re D * Re D)) - Re B * (Re B * ((cmod B)2 - (Re D)2))) / cmod B ^ 4
    using  $\langle$  (cmod B)2 > (cmod A)2  $\rangle$   $\langle$  A = D  $\rangle$ 
    using  $\langle$  is-real D  $\rangle$  cmod-eq-Re[of A]
    by (auto simp add: field-simps)
  also have ... = ((Im B)2 * (Re D)2 - (Re B)2 * ((Re B)2 + (Im B)2 - (Re D)2) / cmod B ^ 4
proof-
  have cmod B * cmod B = Re B * Re B + Im B * Im B
    by (metis cmod-power2 power2-eq-square)
  thus ?thesis
    by (simp add: power2-eq-square)
qed
  also have ... = (((Re D)2 - (Re B)2) * ((Re B)2 + (Im B)2) / cmod B ^ 4
    by (simp add: power2-eq-square field-simps)
  finally have Im-product: Im (i11 / i12) * Im (i21 / i22) = ((Re D)2 - (Re B)2) * ((Re B)2 + (Im B)2) / cmod
B ^ 4
.

show Im (to-complex-cvec i1) * Im (to-complex-cvec i2) < 0  $\longleftrightarrow$  intersects-x-axis-cmat H
proof safe
  assume opposite: Im (to-complex-cvec i1) * Im (to-complex-cvec i2) < 0
  show intersects-x-axis-cmat H
proof-

```

```

have ((Re D)2 - (Re B)2) * ((Re B)2 + (Im B)2) / cmod B ^ 4 < 0
  using Im-product opposite i1 i2
  by simp
hence ((Re D)2 - (Re B)2) * ((Re B)2 + (Im B)2) < 0
  by (simp add: divide-less-0-iff)
hence (Re D)2 < (Re B)2
  by (simp add: mult-less-0-iff not-sum-power2-lt-zero)
thus ?thesis
  using H ⟨A = D⟩ ⟨is-real D⟩
  by auto
qed
next
have *: (∀ k. k * Im B = 1 → k = 0) → Im B = 0
  apply (safe, erule-tac x=1 / Im B in allE)
  using divide-cancel-left by fastforce
assume intersects-x-axis-cmat H
hence Re D = 0 ∨ (Re D)2 < (Re B)2
  using H ⟨A = D⟩
  by auto
hence (Re D)2 < (Re B)2
  using ⟨is-real D⟩ line H ⟨C = cnj B⟩
  using not-x-axis *
  by (auto simp add: complex-eq-iff)
hence ((Re D)2 - (Re B)2) * ((Re B)2 + (Im B)2) < 0
by (metis add-cancel-left-left diff-less-eq mult-eq-0-iff mult-less-0-iff power2-eq-square power2-less-0 sum-squares-gt-zero-iff)
thus Im (to-complex-cvec i1) * Im (to-complex-cvec i2) < 0
  using Im-product i1 i2
  using divide-eq-0-iff divide-less-0-iff prod.simps(2) to-complex-cvec-def zero-complex.simps(1) zero-less-norm-iff
  by fastforce
qed
qed
thus ?thesis
  using assms ⟨calc-ideal-points H = {i1, i2}⟩ ⟨i1 ≠ i2⟩
  by auto
qed

end
theory Poincare-Perpendicular
  imports Poincare-Lines-Axis-Intersections
begin

```

9 H-perpendicular h-lines in the Poincaré model

definition *perpendicular-to-x-axis-cmat* :: *complex-mat* ⇒ *bool* **where**
[simp]: *perpendicular-to-x-axis-cmat* *H* ↔ (let (A, B, C, D) = *H* in *is-real B*)

lift-definition *perpendicular-to-x-axis-clmat* :: *circline-mat* ⇒ *bool* **is** *perpendicular-to-x-axis-cmat*
done

lift-definition *perpendicular-to-x-axis* :: *circline* ⇒ *bool* **is** *perpendicular-to-x-axis-clmat*
by transfer auto

lemma *perpendicular-to-x-axis*:

assumes *is-poincare-line H*
shows *perpendicular-to-x-axis H* ↔ *perpendicular x-axis H*
using *assms*
unfolding *perpendicular-def*

proof (*transfer, transfer*)

fix *H*
assume *hh: hermitean H ∧ H ≠ mat-zero is-poincare-line-cmat H*
obtain *A B C D* **where** *: *H = (A, B, C, D)*
by (*cases H, auto*)
hence *is-real A (cmod B)² > (cmod A)² H = (A, B, cnj B, A)*
using *hermitean-elems[of A B C D] hh*
by auto

thus *perpendicular-to-x-axis-cmat* $H =$
 $(\text{cos-angle-cmat } (\text{of-circline-cmat } x\text{-axis-cmat}) (\text{of-circline-cmat } H) = 0)$
using *cmod-square*[of B] *cmod-square*[of A]
by *simp*
qed

lemma *perpendicular-to-x-axis-y-axis*:
assumes *perpendicular-to-x-axis* (*poincare-line* 0_h (*of-complex* z)) $z \neq 0$
shows *is-imag* z
using *assms*
by (*transfer*, *transfer*, *simp*)

lemma *wlog-perpendicular-axes*:
assumes *in-disc*: $u \in \text{unit-disc } v \in \text{unit-disc } z \in \text{unit-disc}$
assumes *perpendicular*: *is-poincare-line* $H1$ *is-poincare-line* $H2$ *perpendicular* $H1$ $H2$
assumes $z \in \text{circline-set } H1 \cap \text{circline-set } H2$ $u \in \text{circline-set } H1$ $v \in \text{circline-set } H2$
assumes *axes*: $\bigwedge x y. \llbracket \text{is-real } x; 0 \leq \text{Re } x; \text{Re } x < 1; \text{is-imag } y; 0 \leq \text{Im } y; \text{Im } y < 1 \rrbracket \implies P$ 0_h (*of-complex* x)
(*of-complex* y)
assumes *moebius*: $\bigwedge M u v w. \llbracket \text{unit-disc-fix } M; u \in \text{unit-disc}; v \in \text{unit-disc}; w \in \text{unit-disc}; P (\text{moebius-pt } M u)$
(*moebius-pt* $M v$) (*moebius-pt* $M w$) $\rrbracket \implies P$ $u v w$
assumes *conjugate*: $\bigwedge u v w. \llbracket u \in \text{unit-disc}; v \in \text{unit-disc}; w \in \text{unit-disc}; P (\text{conjugate } u) (\text{conjugate } v) (\text{conjugate } w) \rrbracket \implies P$ $u v w$
shows $P z u v$
proof–
have $\forall v H1 H2. \text{is-poincare-line } H1 \wedge \text{is-poincare-line } H2 \wedge \text{perpendicular } H1 H2 \wedge$
 $z \in \text{circline-set } H1 \cap \text{circline-set } H2 \wedge u \in \text{circline-set } H1 \wedge v \in \text{circline-set } H2 \wedge v \in \text{unit-disc} \longrightarrow P$
 $z u v$ (*is ?P z u*)
proof (*rule wlog-x-axis*[*where* $P=?P$])
fix x
assume x : *is-real* x $\text{Re } x \geq 0$ $\text{Re } x < 1$
have *of-complex* $x \in \text{unit-disc}$
using x
by (*simp add: cmod-eq-Re*)
show $?P$ 0_h (*of-complex* x)
proof *safe*
fix $v H1 H2$
assume $v \in \text{unit-disc}$
then obtain y **where** y : $v = \text{of-complex } y$
using *inf-or-of-complex*[of v]
by *auto*

assume 1: *is-poincare-line* $H1$ *is-poincare-line* $H2$ *perpendicular* $H1$ $H2$
assume 2: $0_h \in \text{circline-set } H1$ $0_h \in \text{circline-set } H2$ *of-complex* $x \in \text{circline-set } H1$ $v \in \text{circline-set } H2$

show P 0_h (*of-complex* x) v
proof (*cases of-complex* $x = 0_h$)
case *True*
show P 0_h (*of-complex* x) v
proof (*cases* $v = 0_h$)
case *True*
thus *?thesis*
using $\langle \text{of-complex } x = 0_h \rangle$
using *axes*[of 0 0]
by *simp*
next
case *False*
show *?thesis*
proof (*rule wlog-rotation-to-positive-y-axis*)
show $v \in \text{unit-disc}$ $v \neq 0_h$
by *fact+*
next
fix y
assume *is-imag* y $0 < \text{Im } y$ $\text{Im } y < 1$
thus P 0_h (*of-complex* x) (*of-complex* y)

```

    using  $x$  axes[of  $x$   $y$ ]
    by simp
next
fix  $\varphi$   $u$ 
assume  $u \in \text{unit-disc}$   $u \neq 0_h$ 
     $P$   $0_h$  (of-complex  $x$ ) (moebius-pt (moebius-rotation  $\varphi$ )  $u$ )
thus  $P$   $0_h$  (of-complex  $x$ )  $u$ 
    using  $\langle \text{of-complex } x = 0_h \rangle$ 
    using moebius[of moebius-rotation  $\varphi$   $0_h$   $0_h$   $u$ ]
    by simp
qed
qed
next
case False
hence *: poincare-line  $0_h$  (of-complex  $x$ ) = x-axis
    using x poincare-line-0-real-is-x-axis[of of-complex  $x$ ]
    unfolding circle-set-x-axis
    by auto
hence  $H1 = \text{x-axis}$ 
    using unique-poincare-line[of  $0_h$  of-complex  $x$   $H1$ ] 1 2
    using  $\langle \text{of-complex } x \in \text{unit-disc} \rangle$  False
    by simp
have is-imag  $y$ 
proof (cases  $y = 0$ )
    case True
    thus ?thesis
        by simp
next
case False
hence  $0_h \neq \text{of-complex } y$ 
    using of-complex-zero-iff[of  $y$ ]
    by metis
hence  $H2 = \text{poincare-line } 0_h$  (of-complex  $y$ )
    using 1 2  $\langle v \in \text{unit-disc} \rangle$ 
    using unique-poincare-line[of  $0_h$  of-complex  $y$   $H2$ ]  $y$ 
    by simp
thus ?thesis
    using 1  $\langle H1 = \text{x-axis} \rangle$ 
    using perpendicular-to-x-axis-y-axis[of  $y$ ] False
    using perpendicular-to-x-axis[of  $H2$ ]
    by simp
qed
show  $P$   $0_h$  (of-complex  $x$ )  $v$ 
proof (cases  $\text{Im } y \geq 0$ )
    case True
    thus ?thesis
        using axes[of  $x$   $y$ ]  $x$   $y$   $\langle \text{is-imag } y \rangle$   $\langle v \in \text{unit-disc} \rangle$ 
        by (simp add: cmod-eq-Im)
next
case False
show ?thesis
proof (rule conjugate)
    have  $\text{Im } (\text{cnj } y) < 1$ 
        using  $\langle v \in \text{unit-disc} \rangle$   $y$   $\langle \text{is-imag } y \rangle$  eq-minus-cnj-iff-imag[of  $y$ ]
        by (simp add: cmod-eq-Im)
    thus  $P$  (conjugate  $0_h$ ) (conjugate (of-complex  $x$ )) (conjugate  $v$ )
        using  $\langle \text{is-real } x \rangle$  eq-cnj-iff-real[of  $x$ ]  $y$   $\langle \text{is-imag } y \rangle$ 
        using axes[OF  $x$ , of  $\text{cnj } y$ ] False
        by simp
    show  $0_h \in \text{unit-disc}$  of-complex  $x \in \text{unit-disc}$   $v \in \text{unit-disc}$ 
        by (simp, fact+)
qed
qed
qed
qed
next

```

```

show  $z \in \text{unit-disc } u \in \text{unit-disc}$ 
  by fact+
next
fix  $M u v$ 
assume *:  $\text{unit-disc-fix } M u \in \text{unit-disc } v \in \text{unit-disc}$ 
assume **:  $?P (\text{moebius-pt } M u) (\text{moebius-pt } M v)$ 
show  $?P u v$ 
proof safe
  fix  $w H1 H2$ 
  assume ***:  $\text{is-poincare-line } H1 \text{ is-poincare-line } H2 \text{ perpendicular } H1 H2$ 
     $u \in \text{circline-set } H1 u \in \text{circline-set } H2$ 
     $v \in \text{circline-set } H1 w \in \text{circline-set } H2 w \in \text{unit-disc}$ 
  thus  $P u v w$ 
    using moebius[of  $M u v w$ ] *
    using **[rule-format, of moebius-circline  $M H1$  moebius-circline  $M H2$  moebius-pt  $M w$ ]
    by simp
qed
qed
thus ?thesis
  using assms
  by blast
qed

```

lemma *wlog-perpendicular-foot*:

```

assumes in-disc:  $u \in \text{unit-disc } v \in \text{unit-disc } w \in \text{unit-disc } z \in \text{unit-disc}$ 
assumes perpendicular:  $u \neq v \text{ is-poincare-line } H \text{ perpendicular } (\text{poincare-line } u v) H$ 
assumes  $z \in \text{circline-set } (\text{poincare-line } u v) \cap \text{circline-set } H w \in \text{circline-set } H$ 
assumes axes:  $\bigwedge u v w. \llbracket \text{is-real } u; 0 < \text{Re } u; \text{Re } u < 1; \text{is-real } v; -1 < \text{Re } v; \text{Re } v < 1; \text{Re } u \neq \text{Re } v; \text{is-imag } w; 0 \leq \text{Im } w; \text{Im } w < 1 \rrbracket \implies P 0_h (\text{of-complex } u) (\text{of-complex } v) (\text{of-complex } w)$ 
assumes moebius:  $\bigwedge M z u v w. \llbracket \text{unit-disc-fix } M; u \in \text{unit-disc}; v \in \text{unit-disc}; w \in \text{unit-disc}; z \in \text{unit-disc}; P (\text{moebius-pt } M z) (\text{moebius-pt } M u) (\text{moebius-pt } M v) (\text{moebius-pt } M w) \rrbracket \implies P z u v w$ 
assumes conjugate:  $\bigwedge z u v w. \llbracket u \in \text{unit-disc}; v \in \text{unit-disc}; w \in \text{unit-disc}; P (\text{conjugate } z) (\text{conjugate } u) (\text{conjugate } v) (\text{conjugate } w) \rrbracket \implies P z u v w$ 
assumes perm:  $P z v u w \implies P z u v w$ 
shows  $P z u v w$ 

```

proof—

```

obtain  $m n$  where  $mn: m = u \vee m = v n = u \vee n = v m \neq n m \neq z$ 
  using  $\langle u \neq v \rangle$ 
  by auto

```

```

have  $n \in \text{circline-set } (\text{poincare-line } z m)$ 
  using  $\langle z \in \text{circline-set } (\text{poincare-line } u v) \cap \text{circline-set } H \rangle$ 
  using  $mn$ 
  using unique-poincare-line[of  $z m$  poincare-line  $u v$ , symmetric] in-disc
  by auto

```

have $\forall n. n \in \text{unit-disc} \wedge m \neq n \wedge n \in \text{circline-set } (\text{poincare-line } z m) \wedge m \neq z \longrightarrow P z m n w$ (is ?Q $z m w$)

proof (rule *wlog-perpendicular-axes*[where $P=?Q$])

```

show  $\text{is-poincare-line } (\text{poincare-line } u v)$ 
  using  $\langle u \neq v \rangle$ 
  by auto

```

next

```

show  $\text{is-poincare-line } H$ 
  by fact

```

next

```

show  $m \in \text{unit-disc } m \in \text{circline-set } (\text{poincare-line } u v)$ 
  using  $mn$  in-disc
  by auto

```

next

```

show  $w \in \text{unit-disc } z \in \text{unit-disc}$ 
  by fact+

```

next

```

show  $z \in \text{circline-set } (\text{poincare-line } u v) \cap \text{circline-set } H$ 
  by fact

```

next

```

show  $\text{perpendicular } (\text{poincare-line } u v) H$ 

```

```

  by fact
next
show  $w \in \text{circline-set } H$ 
  by fact
next
fix  $x y$ 
assume  $xy$ :  $\text{is-real } x \ 0 \leq \text{Re } x \ \text{Re } x < 1 \ \text{is-imag } y \ 0 \leq \text{Im } y \ \text{Im } y < 1$ 
show  $?Q \ 0_h \ (\text{of-complex } x) \ (\text{of-complex } y)$ 
proof safe
  fix  $n$ 
  assume  $n \in \text{unit-disc of-complex } x \neq n$ 
  assume  $n \in \text{circline-set (poincare-line } 0_h \ (\text{of-complex } x)) \ \text{of-complex } x \neq 0_h$ 
  hence  $n \in \text{circline-set } x\text{-axis}$ 
  using  $\text{poincare-line-0-real-is-x-axis[of of-complex } x] \ xy$ 
  by  $(\text{auto simp add: circline-set-x-axis})$ 
  then obtain  $n'$  where  $n': n = \text{of-complex } n'$ 
  using  $\text{inf-or-of-complex[of } n] \ \langle n \in \text{unit-disc} \rangle$ 
  by auto
  hence  $\text{is-real } n'$ 
  using  $\langle n \in \text{circline-set } x\text{-axis} \rangle$ 
  using  $\text{of-complex-inj}$ 
  unfolding  $\text{circline-set-x-axis}$ 
  by auto
  hence  $-1 < \text{Re } n' \ \text{Re } n' < 1$ 
  using  $\langle n \in \text{unit-disc} \rangle \ n'$ 
  by  $(\text{auto simp add: cmod-eq-Re})$ 

  have  $\text{Re } n' \neq \text{Re } x$ 
  using  $\text{complex.expand[of } n' \ x] \ \langle \text{is-real } n' \rangle \ \langle \text{is-real } x \rangle \ \langle \text{of-complex } x \neq n \rangle \ n'$ 
  by auto

  have  $\text{Re } x > 0$ 
  using  $xy \ \langle \text{of-complex } x \neq 0_h \rangle$ 
  by  $(\text{cases } \text{Re } x = 0, \ \text{auto simp add: complex.expand})$ 

  show  $P \ 0_h \ (\text{of-complex } x) \ n \ (\text{of-complex } y)$ 
  using  $\text{axes[of } x \ n' \ y] \ xy \ n' \ \langle \text{Re } x > 0 \rangle \ \langle \text{is-real } n' \rangle \ \langle -1 < \text{Re } n' \rangle \ \langle \text{Re } n' < 1 \rangle \ \langle \text{Re } n' \neq \text{Re } x \rangle$ 
  by  $\text{simp}$ 
qed
next
fix  $M \ u \ v \ w$ 
assume 1:  $\text{unit-disc-fix } M \ u \in \text{unit-disc } v \in \text{unit-disc } w \in \text{unit-disc}$ 
assume 2:  $?Q \ (\text{moebius-pt } M \ u) \ (\text{moebius-pt } M \ v) \ (\text{moebius-pt } M \ w)$ 
show  $?Q \ u \ v \ w$ 
proof safe
  fix  $n$ 
  assume  $n \in \text{unit-disc } v \neq n \ n \in \text{circline-set (poincare-line } u \ v) \ v \neq u$ 
  thus  $P \ u \ v \ n \ w$ 
  using  $\text{moebius[of } M \ v \ n \ w \ u] \ 1 \ 2[\text{rule-format, of moebius-pt } M \ n]$ 
  by  $\text{fastforce}$ 
qed
next
fix  $u \ v \ w$ 
assume 1:  $u \in \text{unit-disc } v \in \text{unit-disc } w \in \text{unit-disc}$ 
assume 2:  $?Q \ (\text{conjugate } u) \ (\text{conjugate } v) \ (\text{conjugate } w)$ 
show  $?Q \ u \ v \ w$ 
proof safe
  fix  $n$ 
  assume  $n \in \text{unit-disc } v \neq n \ n \in \text{circline-set (poincare-line } u \ v) \ v \neq u$ 
  thus  $P \ u \ v \ n \ w$ 
  using  $\text{conjugate[of } v \ n \ w \ u] \ 1 \ 2[\text{rule-format, of conjugate } n]$ 
  using  $\text{conjugate-inj}$ 
  by auto
qed
qed
thus  $?thesis$ 

```

```

    using mn in-disc ⟨n ∈ circline-set (poincare-line z m)⟩ perm
  by auto
qed

```

```

lemma perpendicular-to-x-axis-intersects-x-axis:
  assumes is-poincare-line H perpendicular-to-x-axis H
  shows intersects-x-axis H
  using assms hermitean-elems
  by (transfer, transfer, auto simp add: cmod-eq-Re)

```

```

lemma perpendicular-intersects:
  assumes is-poincare-line H1 is-poincare-line H2
  assumes perpendicular H1 H2
  shows ∃ z. z ∈ unit-disc ∧ z ∈ circline-set H1 ∩ circline-set H2 (is ?P' H1 H2)
proof-
  have ∀ H2. is-poincare-line H2 ∧ perpendicular H1 H2 ⟶ ?P' H1 H2 (is ?P H1)
  proof (rule wlog-line-x-axis)
    show ?P x-axis
    proof safe
      fix H2
      assume is-poincare-line H2 perpendicular x-axis H2
      thus ∃ z. z ∈ unit-disc ∧ z ∈ circline-set x-axis ∩ circline-set H2
        using perpendicular-to-x-axis[of H2]
        using perpendicular-to-x-axis-intersects-x-axis[of H2]
        using intersects-x-axis-iff[of H2]
      by auto
    qed
  qed

```

```

next
  fix M
  assume unit-disc-fix M
  assume *: ?P (moebius-circline M H1)
  show ?P H1
  proof safe
    fix H2
    assume is-poincare-line H2 perpendicular H1 H2
    then obtain z where z ∈ unit-disc z ∈ circline-set (moebius-circline M H1) ∧ z ∈ circline-set (moebius-circline
M H2)
      using *[rule-format, of moebius-circline M H2] ⟨unit-disc-fix M⟩
      by auto
    thus ∃ z. z ∈ unit-disc ∧ z ∈ circline-set H1 ∩ circline-set H2
      using ⟨unit-disc-fix M⟩
      by (rule-tac x=moebius-pt (-M) z in exI)
        (metis IntI add.inverse-inverse circline-set-moebius-circline-iff moebius-pt-comp-inv-left uminus-moebius-def
unit-disc-fix-discI unit-disc-fix-moebius-uminus)
    qed
  next
    show is-poincare-line H1
      by fact
  qed
  thus ?thesis
    using assms
    by auto
qed

```

```

definition calc-perpendicular-to-x-axis-cmat :: complex-vec ⇒ complex-mat where
[simp]: calc-perpendicular-to-x-axis-cmat z =
  (let (z1, z2) = z
    in if z1*cnj z2 + z2*cnj z1 = 0 then
      (0, 1, 1, 0)
    else
      let A = z1*cnj z2 + z2*cnj z1;
          B = -(z1*cnj z1 + z2*cnj z2)
      in (A, B, B, A)
  )

```

lift-definition *calc-perpendicular-to-x-axis-clmat* :: *complex-homo-coords* \Rightarrow *circline-mat* **is** *calc-perpendicular-to-x-axis-cmat*
by (*auto simp add: hermitean-def mat-adj-def mat-cnj-def Let-def split: if-split-asm*)

lift-definition *calc-perpendicular-to-x-axis* :: *complex-homo* \Rightarrow *circline* **is** *calc-perpendicular-to-x-axis-clmat*

proof (*transfer*)

fix $z w$

assume $z \neq \text{vec-zero } w \neq \text{vec-zero}$

obtain $z1 z2 w1 w2$ **where** $zw: z = (z1, z2) w = (w1, w2)$

by (*cases z, cases w, auto*)

assume $z \approx_v w$

then obtain k **where** $*: k \neq 0 w1 = k*z1 w2 = k*z2$

using zw

by *auto*

have $w1 * \text{cnj } w2 + w2 * \text{cnj } w1 = (k * \text{cnj } k) * (z1 * \text{cnj } z2 + z2 * \text{cnj } z1)$

using $*$

by (*auto simp add: field-simps*)

moreover

have $w1 * \text{cnj } w1 + w2 * \text{cnj } w2 = (k * \text{cnj } k) * (z1 * \text{cnj } z1 + z2 * \text{cnj } z2)$

using $*$

by (*auto simp add: field-simps*)

ultimately

show *circline-eq-cmat* (*calc-perpendicular-to-x-axis-cmat* z) (*calc-perpendicular-to-x-axis-cmat* w)

using $zw *$

apply (*auto simp add: Let-def*)

apply (*rule-tac x=Re (k * cnj k) in exI, auto simp add: complex.expand field-simps*)

done

qed

lemma *calc-perpendicular-to-x-axis*:

assumes $z \neq \text{of-complex } 1 z \neq \text{of-complex } (-1)$

shows $z \in \text{circline-set } (\text{calc-perpendicular-to-x-axis } z) \wedge$

$\text{is-poincare-line } (\text{calc-perpendicular-to-x-axis } z) \wedge$

$\text{perpendicular-to-x-axis } (\text{calc-perpendicular-to-x-axis } z)$

using *assms*

unfolding *circline-set-def perpendicular-def*

proof (*simp, transfer, transfer*)

fix $z :: \text{complex-vec}$

obtain $z1 z2$ **where** $z: z = (z1, z2)$

by (*cases z, auto*)

assume $** : \neg z \approx_v \text{of-complex-cvec } 1 \neg z \approx_v \text{of-complex-cvec } (-1)$

show $\text{on-circline-cmat-cvec } (\text{calc-perpendicular-to-x-axis-cmat } z) z \wedge$

$\text{is-poincare-line-cmat } (\text{calc-perpendicular-to-x-axis-cmat } z) \wedge$

$\text{perpendicular-to-x-axis-cmat } (\text{calc-perpendicular-to-x-axis-cmat } z)$

proof (*cases z1*cnj z2 + z2*cnj z1 = 0*)

case *True*

thus *?thesis*

using z

by (*simp add: vec-cnj-def hermitean-def mat-adj-def mat-cnj-def mult.commute*)

next

case *False*

hence $z2 \neq 0$

using z

by *auto*

hence $\text{Re } (z2 * \text{cnj } z2) \neq 0$

using $\langle z2 \neq 0 \rangle$

by (*auto simp add: complex.expand*)

have $z1 \neq -z2 \wedge z1 \neq z2$

proof (*rule ccontr*)

assume $\neg ?thesis$

hence $z \approx_v \text{of-complex-cvec } 1 \vee z \approx_v \text{of-complex-cvec } (-1)$

using $z \langle z2 \neq 0 \rangle$

by *auto*

thus *False*

using $**$


```

    by auto
qed

let ?A = z1*cnj z2 + z2*cnj z1 and ?B = -(z1*cnj z1 + z2*cnj z2)
have Re(z1*cnj z1 + z2*cnj z2) ≥ 0
  by auto
hence Re ?B ≤ 0
  by (smt uminus-complex.simps(1))
hence abs (Re ?B) = - Re ?B
  by auto
also have ... = (Re z1)2 + (Im z1)2 + (Re z2)2 + (Im z2)2
  by (simp add: power2-eq-square[symmetric])
also have ... > abs (Re ?A)
proof (cases Re ?A ≥ 0)
  case False
  have (Re z1 + Re z2)2 + (Im z1 + Im z2)2 > 0
    using ⟨z1 ≠ -z2 ∧ z1 ≠ z2⟩
    by (metis add.commute add.inverse-unique complex-neq-0 plus-complex.code plus-complex.simps)
  thus ?thesis
    using False
    by (simp add: power2-sum power2-eq-square field-simps)
next
  case True
  have (Re z1 - Re z2)2 + (Im z1 - Im z2)2 > 0
    using ⟨z1 ≠ -z2 ∧ z1 ≠ z2⟩
    by (meson complex-eq-iff right-minus-eq sum-power2-gt-zero-iff)
  thus ?thesis
    using True
    by (simp add: power2-sum power2-eq-square field-simps)
qed
finally
have abs (Re ?B) > abs (Re ?A)
.
moreover
have cmod ?B = abs (Re ?B) cmod ?A = abs (Re ?A)
  by (simp-all add: cmod-eq-Re)
ultimately
have (cmod ?B)2 > (cmod ?A)2
  by (smt power2-le-imp-le)
thus ?thesis
  using z False
  by (simp-all add: Let-def hermitean-def mat-adj-def mat-cnj-def cmod-eq-Re vec-cnj-def field-simps)
qed
qed

```

lemma *ex-perpendicular*:

```

assumes is-poincare-line H z ∈ unit-disc
shows ∃ H'. is-poincare-line H' ∧ perpendicular H H' ∧ z ∈ circline-set H' (is ?P' H z)
proof-
have ∀ z. z ∈ unit-disc → ?P' H z (is ?P H)
proof (rule wlog-line-x-axis)
show ?P x-axis
proof safe
fix z
assume z ∈ unit-disc
then have z ≠ of-complex 1 z ≠ of-complex (-1)
  by auto
thus ?P' x-axis z
  using ⟨z ∈ unit-disc⟩
  using calc-perpendicular-to-x-axis[of z] perpendicular-to-x-axis
  by (rule-tac x = calc-perpendicular-to-x-axis z in exI, auto)
qed
next
fix M
assume unit-disc-fix M
assume *: ?P (moebius-circline M H)

```

```

show ?P H
proof safe
  fix z
  assume z ∈ unit-disc
  hence moebius-pt M z ∈ unit-disc
    using ⟨unit-disc-fix M⟩
  by auto
then obtain H' where *: is-poincare-line H' perpendicular (moebius-circline M H) H' moebius-pt M z ∈ circline-set
H'
  using *
  by auto
have h: H = moebius-circline (-M) (moebius-circline M H)
  by auto
show ?P' H z
  using * ⟨unit-disc-fix M⟩
  apply (subst h)
  apply (rule-tac x=moebius-circline (-M) H' in exI)
  apply (simp del: moebius-circline-comp-inv-left)
  done
qed
qed fact
thus ?thesis
  using assms
  by simp
qed

lemma ex-perpendicular-foot:
  assumes is-poincare-line H z ∈ unit-disc
  shows ∃ H'. is-poincare-line H' ∧ z ∈ circline-set H' ∧ perpendicular H H' ∧
    (∃ z' ∈ unit-disc. z' ∈ circline-set H' ∩ circline-set H)
  using assms
  using ex-perpendicular[OF assms]
  using perpendicular-intersects[of H]
  by blast

lemma Pythagoras:
  assumes in-disc: u ∈ unit-disc v ∈ unit-disc w ∈ unit-disc v ≠ w
  assumes distinct[u, v, w] → perpendicular (poincare-line u v) (poincare-line u w)
  shows cosh (poincare-distance v w) = cosh (poincare-distance u v) * cosh (poincare-distance u w) (is ?P' u v w)
proof (cases distinct [u, v, w])
  case False
  thus ?thesis
    using in-disc
    by (auto simp add: poincare-distance-sym)
next
  case True
  have distinct [u, v, w] → ?P' u v w (is ?P u v w)
  proof (rule wlog-perpendicular-axes[where P=?P])
    show is-poincare-line (poincare-line u v) is-poincare-line (poincare-line u w)
      using ⟨distinct [u, v, w]⟩
      by simp-all
  next
    show perpendicular (poincare-line u v) (poincare-line u w)
      using True assms
      by simp
  next
    show u ∈ unit-disc v ∈ unit-disc w ∈ unit-disc
      by fact+
  next
    show v ∈ circline-set (poincare-line u v) w ∈ circline-set (poincare-line u w)
      u ∈ circline-set (poincare-line u v) ∩ circline-set (poincare-line u w)
      using ⟨distinct [u, v, w]⟩
      by auto
  next
    fix x y
    assume x: is-real x 0 ≤ Re x Re x < 1

```

```

assume y: is-imag y 0 ≤ Im y Im y < 1

have of-complex x ∈ unit-disc of-complex y ∈ unit-disc
  using x y
  by (simp-all add: cmod-eq-Re cmod-eq-Im)

show ?P 0h (of-complex x) (of-complex y)
proof
  assume distinct [0h, of-complex x, of-complex y]
  hence x ≠ 0 y ≠ 0
    by auto

  let ?den1 = 1 - (cmod x)2 and ?den2 = 1 - (cmod y)2
  have ?den1 > 0 ?den2 > 0
    using x y
    by (simp-all add: cmod-eq-Re cmod-eq-Im abs-square-less-1)

  let ?d1 = 1 + 2 * (cmod x)2 / ?den1
  have cosh (poincare-distance 0h (of-complex x)) = ?d1
    using ⟨?den1 > 0⟩
    using poincare-distance-formula[of 0h of-complex x] ⟨of-complex x ∈ unit-disc⟩
    by simp

  moreover

  let ?d2 = 1 + 2 * (cmod y)2 / ?den2
  have cosh (poincare-distance 0h (of-complex y)) = ?d2
    using ⟨?den2 > 0⟩ ⟨of-complex y ∈ unit-disc⟩
    using poincare-distance-formula[of 0h of-complex y]
    by simp

  moreover

  let ?den = ?den1 * ?den2
  let ?d3 = 1 + 2 * (cmod (x - y))2 / ?den
  have cosh (poincare-distance (of-complex x) (of-complex y)) = ?d3
    using ⟨of-complex x ∈ unit-disc⟩ ⟨of-complex y ∈ unit-disc⟩
    using ⟨?den1 > 0⟩ ⟨?den2 > 0⟩
    using poincare-distance-formula[of of-complex x of-complex y]
    by simp

  moreover
  have ?d1 * ?d2 = ?d3
  proof-
    have ?d3 = ((1 - (cmod x)2) * (1 - (cmod y)2) + 2 * (cmod (x - y))2) / ?den
      using ⟨?den1 > 0⟩ ⟨?den2 > 0⟩
      by (subst add-num-frac, simp, simp)
    also have ... = (Re ((1 - x * cnj x) * (1 - y * cnj y) + 2 * (x - y) * cnj (x - y))) / ?den
      using ⟨is-real x⟩ ⟨is-imag y⟩
      by ((subst cmod-square)+, simp)
    also have ... = Re (1 + x * cnj x * y * cnj y
      + x * cnj x - 2 * y * cnj x - 2 * x * cnj y + y * cnj y) / ?den
      by (simp add: field-simps)
    also have ... = Re ((1 + y * cnj y) * (1 + x * cnj x)) / ?den
      using ⟨is-real x⟩ ⟨is-imag y⟩
      by (simp add: field-simps)
  finally
  show ?thesis
    using ⟨?den1 > 0⟩ ⟨?den2 > 0⟩
    apply (subst add-num-frac, simp)
    apply (subst add-num-frac, simp)
    apply simp
    apply (subst cmod-square)+
    apply (simp add: field-simps)
  done

qed
ultimately
show ?P' 0h (of-complex x) (of-complex y)

```

```

    by simp
qed
next
fix M u v w
assume 1: unit-disc-fix M u ∈ unit-disc v ∈ unit-disc w ∈ unit-disc
assume 2: ?P (moebius-pt M u) (moebius-pt M v) (moebius-pt M w)
show ?P u v w
  using 1 2
  by auto
next
fix u v w
assume 1: u ∈ unit-disc v ∈ unit-disc w ∈ unit-disc
assume 2: ?P (conjugate u) (conjugate v) (conjugate w)
show ?P u v w
  using 1 2
  by (auto simp add: conjugate-inj)
qed
thus ?thesis
  using True
  by simp
qed
end

```

10 Poincaré disc model types

In this section we introduce datatypes that represent objects in the Poincaré disc model. The types are defined as subtypes (e.g., the h-points are defined as elements of $\mathbb{C}P^1$ that lie within the unit disc). The functions on those types are defined by lifting the functions defined on the carrier type (e.g., h-distance is defined by lifting the distance function defined for elements of $\mathbb{C}P^1$).

```

theory Poincare
imports Poincare-Lines Poincare-Between Poincare-Distance Poincare-Circles
begin

```

10.1 H-points

```

typedef p-point = {z. z ∈ unit-disc}
  using zero-in-unit-disc
  by (rule-tac x=0h in exI, simp)

```

```

setup-lifting type-definition-p-point

```

Point zero

```

lift-definition p-zero :: p-point is 0h
  by (rule zero-in-unit-disc)

```

Constructing h-points from complex numbers

```

lift-definition p-of-complex :: complex ⇒ p-point is λ z. if cmod z < 1 then of-complex z else 0h
  by auto

```

10.2 H-lines

```

typedef p-line = {H. is-poincare-line H}
  by (rule-tac x=x-axis in exI, simp)

```

```

setup-lifting type-definition-p-line

```

```

lift-definition p-incident :: p-line ⇒ p-point ⇒ bool is on-circline
  done

```

Set of h-points on an h-line

```

definition p-points :: p-line ⇒ p-point set where
  p-points l = {p. p-incident l p}

```

x-axis is an example of an h-line

```
lift-definition p-x-axis :: p-line is x-axis
  by simp
```

Constructing the unique h-line from two h-points

```
lift-definition p-line :: p-point  $\Rightarrow$  p-point  $\Rightarrow$  p-line is poincare-line
proof-
  fix u v
  show is-poincare-line (poincare-line u v)
  proof (cases u  $\neq$  v)
    case True
    thus ?thesis
    by simp
  next
```

This branch must work only for formal reasons.

```
    case False
    thus ?thesis
    by (transfer, transfer, auto simp add: hermitean-def mat-adj-def mat-cnj-def split: if-split-asm)
  qed
qed
```

Next we show how to lift some lemmas. This could be done for all the lemmas that we have proved earlier, but we do not do that.

If points are different then the constructed line contains the starting points

```
lemma p-on-line:
  assumes z  $\neq$  w
  shows p-incident (p-line z w) z
    p-incident (p-line z w) w
  using assms
  by (transfer, simp)+
```

There is a unique h-line passing through the two different given h-points

```
lemma
  assumes u  $\neq$  v
  shows  $\exists! l. \{u, v\} \subseteq p\text{-points } l$ 
  using assms
  apply (rule-tac a=p-line u v in exII, auto simp add: p-points-def p-on-line)
  apply (transfer, simp add: unique-poincare-line)
  done
```

The unique h-line through zero and a non-zero h-point on the x-axis is the x-axis

```
lemma
  assumes p-zero  $\in$  p-points l u  $\in$  p-points l u  $\neq$  p-zero u  $\in$  p-points p-x-axis
  shows l = p-x-axis
  using assms
  unfolding p-points-def
  apply simp
  apply transfer
  using is-poincare-line-0-real-is-x-axis inf-notin-unit-disc
  unfolding circline-set-def
  by blast
```

10.3 H-collinearity

```
lift-definition p-collinear :: p-point set  $\Rightarrow$  bool is poincare-collinear
  done
```

10.4 H-isometries

H-isometries are functions that map the unit disc onto itself

```
typedef p-isometry = {f. unit-disc-fix f}
```

by (rule-tac x=id in exI, simp add: unit-disc-fix-f-def, rule-tac x=id-moebius in exI, simp)

setup-lifting type-definition-p-isometry

Action of an h-isometry on an h-point

lift-definition *p-isometry-pt* :: *p-isometry* \Rightarrow *p-point* \Rightarrow *p-point* **is** $\lambda f p. f p$
using *unit-disc-fix-f-unit-disc*
by *auto*

Action of an h-isometry on an h-line

lift-definition *p-isometry-line* :: *p-isometry* \Rightarrow *p-line* \Rightarrow *p-line* **is** $\lambda f l. \text{unit-disc-fix-f-circline } f l$
proof–
fix *f H*
assume *unit-disc-fix-f f is-poincare-line H*
then obtain *M* **where** *unit-disc-fix M* **and** $*$: $f = \text{moebius-pt } M \vee f = \text{moebius-pt } M \circ \text{conjugate}$
unfolding *unit-disc-fix-f-def*
by *auto*
show *is-poincare-line (unit-disc-fix-f-circline f H)*
using $*$
proof
assume $f = \text{moebius-pt } M$
thus *?thesis*
using $\langle \text{unit-disc-fix } M \rangle \langle \text{is-poincare-line } H \rangle$
using *unit-disc-fix-f-circline-direct[of M f H]*
by *auto*
next
assume $f = \text{moebius-pt } M \circ \text{conjugate}$
thus *?thesis*
using $\langle \text{unit-disc-fix } M \rangle \langle \text{is-poincare-line } H \rangle$
using *unit-disc-fix-f-circline-indirect[of M f H]*
by *auto*
qed
qed

An example lemma about h-isometries.

H-isometries preserve h-collinearity

lemma *p-collinear-p-isometry-pt* [*simp*]:
shows *p-collinear (p-isometry-pt M ' A)* \longleftrightarrow *p-collinear A*
proof–
have $*$: $\forall M A. ((\lambda x. \text{moebius-pt } M (\text{conjugate } x)) ' A = \text{moebius-pt } M ' (\text{conjugate } ' A))$
by *auto*
show *?thesis*
by *transfer (auto simp add: unit-disc-fix-f-def *)*
qed

10.5 H-distance and h-congruence

lift-definition *p-dist* :: *p-point* \Rightarrow *p-point* \Rightarrow *real* **is** *poincare-distance*
done

definition *p-congruent* :: *p-point* \Rightarrow *p-point* \Rightarrow *p-point* \Rightarrow *p-point* \Rightarrow *bool* **where**
[*simp*]: *p-congruent u v u' v'* \longleftrightarrow *p-dist u v = p-dist u' v'*

lemma

assumes *p-dist u v = p-dist u' v'*
assumes *p-dist v w = p-dist v' w'*
assumes *p-dist u w = p-dist u' w'*
shows $\exists f. \text{p-isometry-pt } f u = u' \wedge \text{p-isometry-pt } f v = v' \wedge \text{p-isometry-pt } f w = w'$
using *assms*
apply *transfer*
using *unit-disc-fix-f-congruent-triangles*
by *auto*

We prove that unit disc equipped with Poincaré distance is a metric space, i.e. an instantiation of *metric-space locale*.

```

instantiation p-point :: metric-space
begin
definition dist-p-point = p-dist
definition (uniformity-p-point :: (p-point × p-point) filter) = (INF e ∈ {0 < ..}. principal {(x, y). dist-class.dist x y < e})
definition open-p-point (U :: p-point set) = (∀ x ∈ U. eventually (λ(x', y). x' = x → y ∈ U) uniformity)
instance
proof
  fix x y :: p-point
  show (dist-class.dist x y = 0) = (x = y)
    unfolding dist-p-point-def
    by (transfer, simp add: poincare-distance-eq-0-iff)
next
  fix x y z :: p-point
  show dist-class.dist x y ≤ dist-class.dist x z + dist-class.dist y z
    unfolding dist-p-point-def
    apply transfer
    using poincare-distance-triangle-inequality poincare-distance-sym
    by metis
qed (simp-all add: open-p-point-def uniformity-p-point-def)
end

```

10.6 H-betweenness

```

lift-definition p-between :: p-point ⇒ p-point ⇒ p-point ⇒ bool is poincare-between
  done
end

```

11 Poincaré model satisfies Tarski axioms

```

theory Poincare-Tarski
  imports Poincare Poincare-Lines-Axis-Intersections Tarski
begin

```

11.1 Pasch axiom

```

lemma Pasch-fun-mono:
  fixes r1 r2 :: real
  assumes  $0 < r1$  and  $r1 \leq r2$  and  $r2 < 1$ 
  shows  $r1 + 1/r1 \geq r2 + 1/r2$ 
proof (cases  $r1 = r2$ )
  case True
  thus ?thesis
    by simp
next
  case False
  hence  $r2 - r1 > 0$ 
    using assms
    by simp

  have  $r1 * r2 < 1$ 
    using assms
    by (smt mult-le-cancel-left1)
  hence  $1 / (r1 * r2) > 1$ 
    using assms
    by simp
  hence  $(r2 - r1) / (r1 * r2) > (r2 - r1)$ 
    using  $\langle r2 - r1 > 0 \rangle$ 
    using mult-less-cancel-left-pos[of  $r2 - r1$   $1$   $1 / (r1 * r2)$ ]
    by simp
  hence  $1 / r1 - 1 / r2 > r2 - r1$ 
    using assms
    by (simp add: field-simps)
  thus ?thesis
    by simp

```

qed

Pasch axiom, non-degenerative case.

lemma *Pasch-nondeg*:

assumes $x \in \text{unit-disc}$ **and** $y \in \text{unit-disc}$ **and** $z \in \text{unit-disc}$ **and** $u \in \text{unit-disc}$ **and** $v \in \text{unit-disc}$

assumes $\text{distinct } [x, y, z, u, v]$

assumes $\neg \text{poincare-collinear } \{x, y, z\}$

assumes $\text{poincare-between } x \ u \ z$ **and** $\text{poincare-between } y \ v \ z$

shows $\exists a. a \in \text{unit-disc} \wedge \text{poincare-between } u \ a \ y \wedge \text{poincare-between } x \ a \ v$

proof –

have $\forall y \ z \ u. \text{distinct } [x, y, z, u, v] \wedge \neg \text{poincare-collinear } \{x, y, z\} \wedge y \in \text{unit-disc} \wedge z \in \text{unit-disc} \wedge u \in \text{unit-disc} \wedge \text{poincare-between } x \ u \ z \wedge \text{poincare-between } y \ v \ z \longrightarrow (\exists a. a \in \text{unit-disc} \wedge \text{poincare-between } u \ a \ y \wedge \text{poincare-between } x \ a \ v)$ (**is** $?P \ x \ v$)

proof (*rule wlog-positive-x-axis*[**where** $P = ?P$])

fix v

assume $v: \text{is-real } v \ 0 < \text{Re } v \ \text{Re } v < 1$

hence $\text{of-complex } v \in \text{unit-disc}$

by (*auto simp add: cmod-eq-Re*)

show $?P \ 0_h$ (*of-complex* v)

proof *safe*

fix $y \ z \ u$

assume $\text{distinct}: \text{distinct } [0_h, y, z, u, \text{of-complex } v]$

assume $\text{in-disc}: y \in \text{unit-disc} \ z \in \text{unit-disc} \ u \in \text{unit-disc}$

then obtain $y' \ z' \ u'$

where $*$: $y = \text{of-complex } y'$ $z = \text{of-complex } z'$ $u = \text{of-complex } u'$

using $\text{inf-or-of-complex inf-notin-unit-disc}$

by *metis*

have $y' \neq 0 \ z' \neq 0 \ u' \neq 0 \ v \neq 0 \ y' \neq z' \ y' \neq u' \ z' \neq u' \ y \neq z \ y \neq u \ z \neq u$

using $\text{of-complex-inj distinct } *$

by *auto*

note $\text{distinct} = \text{distinct this}$

assume $\neg \text{poincare-collinear } \{0_h, y, z\}$

hence $\text{nondeg-yz}: y' * \text{cnj } z' \neq \text{cnj } y' * z'$

using $* \text{poincare-collinear-zero-iff}[of \ y' \ z'] \text{ in-disc distinct}$

by *auto*

assume $\text{poincare-between } 0_h \ u \ z$

hence $\text{Arg } u' = \text{Arg } z' \ \text{cmod } u' \leq \text{cmod } z'$

using $* \text{poincare-between-0uw}[of \ u \ z] \text{ distinct in-disc}$

by *auto*

then obtain $\varphi \ r_u \ r_z$ **where**

$uz\text{-polar}: u' = \text{cor } r_u * \text{cis } \varphi \ z' = \text{cor } r_z * \text{cis } \varphi \ 0 < r_u \ r_u \leq r_z \ 0 < r_z$ **and**

$\varphi = \text{Arg } u' \ \varphi = \text{Arg } z'$

using $* \langle u' \neq 0 \rangle \langle z' \neq 0 \rangle$

by (*smt cmod-cis norm-le-zero-iff*)

obtain $\vartheta \ r_y$ **where**

$y\text{-polar}: y' = \text{cor } r_y * \text{cis } \vartheta \ r_y > 0$ **and** $\vartheta = \text{Arg } y'$

using $\langle y' \neq 0 \rangle$

by (*smt cmod-cis norm-le-zero-iff*)

from $\text{in-disc} * \langle u' = \text{cor } r_u * \text{cis } \varphi \rangle \langle z' = \text{cor } r_z * \text{cis } \varphi \rangle \langle y' = \text{cor } r_y * \text{cis } \vartheta \rangle$

have $r_u < 1 \ r_z < 1 \ r_y < 1$

by (*auto simp: norm-mult*)

note $\text{polar} = \text{this } y\text{-polar } uz\text{-polar}$

have $\text{nondeg}: \text{cis } \vartheta * \text{cis } (-\varphi) \neq \text{cis } (-\vartheta) * \text{cis } \varphi$

using nondeg-yz polar

by *simp*


```

let ?yz = poincare-line y z
let ?v = calc-x-axis-intersection ?yz

assume poincare-between y (of-complex v) z

hence of-complex v ∈ circline-set ?yz
  using in-disc ⟨of-complex v ∈ unit-disc⟩
  using distinct poincare-between-poincare-collinear[of y of-complex v z]
  using unique-poincare-line[of y z]
  by (auto simp add: poincare-collinear-def)
moreover
have of-complex v ∈ circline-set x-axis
  using ⟨is-real v⟩
  unfolding circline-set-x-axis
  by auto
moreover
have ?yz ≠ x-axis
proof (rule ccontr)
  assume ¬ ?thesis
  hence {0h, y, z} ⊆ circline-set (poincare-line y z)
    unfolding circline-set-def
    using distinct poincare-line[of y z]
    by auto
  hence poincare-collinear {0h, y, z}
    unfolding poincare-collinear-def
    using distinct
    by force
  thus False
    using ⟨¬ poincare-collinear {0h, y, z}⟩
    by simp
qed
ultimately
have ?v = of-complex v intersects-x-axis ?yz
  using unique-calc-x-axis-intersection[of poincare-line y z of-complex v]
  using intersects-x-axis-iff[of ?yz]
  using distinct ⟨of-complex v ∈ unit-disc⟩
  by (metis IntI is-poincare-line-poincare-line)+

have intersects-x-axis-positive ?yz
  using ⟨Re v > 0⟩ ⟨of-complex v ∈ unit-disc⟩
  using ⟨of-complex v ∈ circline-set ?yz⟩ ⟨of-complex v ∈ circline-set x-axis⟩
  using intersects-x-axis-positive-iff[of ?yz] ⟨y ≠ z⟩ ⟨?yz ≠ x-axis⟩
  unfolding positive-x-axis-def
  by force

have y ∉ circline-set x-axis
proof (rule ccontr)
  assume ¬ ?thesis
  moreover
  hence poincare-line y (of-complex v) = x-axis
    using distinct ⟨of-complex v ∈ circline-set x-axis⟩
    using in-disc ⟨of-complex v ∈ unit-disc⟩
    using unique-poincare-line[of y of-complex v x-axis]
    by simp
  moreover
  have z ∈ circline-set (poincare-line y (of-complex v))
    using ⟨of-complex v ∈ circline-set ?yz⟩
    using unique-poincare-line[of y of-complex v poincare-line y z]
    using in-disc ⟨of-complex v ∈ unit-disc⟩ distinct
    using poincare-line[of y z]
    unfolding circline-set-def
    by (metis distinct-length-2-or-more is-poincare-line-poincare-line mem-Collect-eq)
  ultimately
  have y ∈ circline-set x-axis z ∈ circline-set x-axis
    by auto

```

hence *poincare-collinear* $\{0_h, y, z\}$
unfolding *poincare-collinear-def*
by *force*
thus *False*
using $\langle \neg \text{poincare-collinear } \{0_h, y, z\} \rangle$
by *simp*
qed

moreover

have $z \notin \text{circline-set } x\text{-axis}$

proof (*rule ccontr*)

assume $\neg ?thesis$

moreover

hence *poincare-line* z (*of-complex* v) = *x-axis*

using *distinct* $\langle \text{of-complex } v \in \text{circline-set } x\text{-axis} \rangle$

using *in-disc* $\langle \text{of-complex } v \in \text{unit-disc} \rangle$

using *unique-poincare-line*[*of* z *of-complex* v *x-axis*]

by *simp*

moreover

have $y \in \text{circline-set } (\text{poincare-line } z \text{ (of-complex } v))$

using $\langle \text{of-complex } v \in \text{circline-set } ?yz \rangle$

using *unique-poincare-line*[*of* z *of-complex* v *poincare-line* y z]

using *in-disc* $\langle \text{of-complex } v \in \text{unit-disc} \rangle$ *distinct*

using *poincare-line*[*of* y z]

unfolding *circline-set-def*

by (*metis distinct-length-2-or-more is-poincare-line-poincare-line mem-Collect-eq*)

ultimately

have $y \in \text{circline-set } x\text{-axis}$ $z \in \text{circline-set } x\text{-axis}$

by *auto*

hence *poincare-collinear* $\{0_h, y, z\}$

unfolding *poincare-collinear-def*

by *force*

thus *False*

using $\langle \neg \text{poincare-collinear } \{0_h, y, z\} \rangle$

by *simp*

qed

ultimately

have $\varphi * \vartheta < 0$

using $\langle \text{poincare-between } y \text{ (of-complex } v) \ z \rangle$

using *poincare-between-x-axis-intersection*[*of* y z *of-complex* v]

using *in-disc* $\langle \text{of-complex } v \in \text{unit-disc} \rangle$ *distinct*

using $\langle \text{of-complex } v \in \text{circline-set } ?yz \rangle$ $\langle \text{of-complex } v \in \text{circline-set } x\text{-axis} \rangle$

using $\langle \varphi = \text{Arg } z' \rangle$ $\langle \vartheta = \text{Arg } y' \rangle *$

by (*simp add: field-simps*)

have $\varphi \neq \pi$ $\varphi \neq 0$

using $\langle z \notin \text{circline-set } x\text{-axis} \rangle *$ *polar cis-pi*

unfolding *circline-set-x-axis*

by *auto*

have $\vartheta \neq \pi$ $\vartheta \neq 0$

using $\langle y \notin \text{circline-set } x\text{-axis} \rangle *$ *polar cis-pi*

unfolding *circline-set-x-axis*

by *auto*

have *phi-sin*: $\varphi > 0 \iff \sin \varphi > 0$ $\varphi < 0 \iff \sin \varphi < 0$

using $\langle \varphi = \text{Arg } z' \rangle$ $\langle \varphi \neq 0 \rangle$ $\langle \varphi \neq \pi \rangle$

using *Arg-bounded*[*of* z']

by (*smt sin-gt-zero sin-le-zero sin-pi-minus sin-0-iff-canon sin-ge-zero*) $+$

have *theta-sin*: $\vartheta > 0 \iff \sin \vartheta > 0$ $\vartheta < 0 \iff \sin \vartheta < 0$

using $\langle \vartheta = \text{Arg } y' \rangle$ $\langle \vartheta \neq 0 \rangle$ $\langle \vartheta \neq \pi \rangle$

using *Arg-bounded*[*of* y']

```

by (smt sin-gt-zero sin-le-zero sin-pi-minus sin-0-iff-canon sin-ge-zero)+

have sin  $\varphi * \sin \vartheta < 0$ 
  using  $\langle \varphi * \vartheta < 0 \rangle$  phi-sin theta-sin
  by (simp add: mult-less-0-iff)

have sin  $(\varphi - \vartheta) \neq 0$ 
proof (rule ccontr)
  assume  $\neg ?thesis$ 
  hence sin  $(\varphi - \vartheta) = 0$ 
    by simp
  have  $-2 * \pi < \varphi - \vartheta < 2 * \pi$ 
    using  $\langle \varphi = \text{Arg } z' \rangle \langle \vartheta = \text{Arg } y' \rangle$  Arg-bounded[of z'] Arg-bounded[of y']  $\langle \varphi \neq \pi \rangle \langle \vartheta \neq \pi \rangle$ 
    by auto
  hence  $\varphi - \vartheta = -\pi \vee \varphi - \vartheta = 0 \vee \varphi - \vartheta = \pi$ 
    using  $\langle \sin(\varphi - \vartheta) = 0 \rangle$ 
    by (smt sin-0-iff-canon sin-periodic-pi2)
  moreover
  {
    assume  $\varphi - \vartheta = -\pi$ 
    hence  $\varphi = \vartheta - \pi$ 
      by simp
    hence False
      using nondeg-yz
      using  $\langle y' = \text{cor } ry * \text{cis } \vartheta \rangle \langle z' = \text{cor } rz * \text{cis } \varphi \rangle \langle rz > 0 \rangle \langle ry > 0 \rangle$ 
      by auto
  }
  moreover
  {
    assume  $\varphi - \vartheta = 0$ 
    hence  $\varphi = \vartheta$ 
      by simp
    hence False
      using  $\langle y' = \text{cor } ry * \text{cis } \vartheta \rangle \langle z' = \text{cor } rz * \text{cis } \varphi \rangle \langle rz > 0 \rangle \langle ry > 0 \rangle$ 
      using nondeg-yz
      by auto
  }
  moreover
  {
    assume  $\varphi - \vartheta = \pi$ 
    hence  $\varphi = \vartheta + \pi$ 
      by simp
    hence False
      using  $\langle y' = \text{cor } ry * \text{cis } \vartheta \rangle \langle z' = \text{cor } rz * \text{cis } \varphi \rangle \langle rz > 0 \rangle \langle ry > 0 \rangle$ 
      using nondeg-yz
      by auto
  }
  ultimately
  show False
    by auto
qed

have  $u \notin \text{circline-set } x\text{-axis}$ 
proof-
  have  $\neg \text{is-real } u'$ 
    using * polar in-disc
    using  $\langle \varphi \neq 0 \rangle \langle \varphi = \text{Arg } u' \rangle \langle \varphi \neq \pi \rangle$  phi-sin(1) phi-sin(2)
    by (metis is-real-arg2)
  moreover
  have  $u \neq \infty_h$ 
    using in-disc
    by auto
  ultimately
  show ?thesis
    using * of-complex-inj[of u']
    unfolding circline-set-x-axis

```

```

    by auto
qed

let ?yu = poincare-line y u
have nondeg-yu: y' * cnj u' ≠ cnj u' * u'
  using nondeg-yz polar ⟨ru > 0⟩ ⟨rz > 0⟩ distinct
  by auto

{
  fix r :: real
  assume r > 0

  have den: cor ry * cis ϑ * cnj 1 * cnj (cor r * cis ϕ) * 1 - cor r * cis ϕ * cnj 1 * cnj (cor ry * cis ϑ) * 1 ≠ 0
    using ⟨0 < r⟩ ⟨0 < ry⟩ nondeg
    by auto

  let ?A = 2 * r * ry * sin(ϕ - ϑ)
  let ?B = i * (r * cis ϕ * (1 + ry2) - ry * cis ϑ * (1 + r2))
  let ?ReB = ry * (1 + r2) * sin ϑ - r * (1 + ry2) * sin ϕ

  have Re (i * (r * cis (-ϕ) * ry * cis (ϑ) - ry * cis (-ϑ) * r * cis (ϕ))) = ?A
    by (simp add: sin-diff field-simps)
  moreover
  have cor ry * cis (-ϑ) * (cor ry * cis ϑ) = ry2 cor r * cis (-ϕ) * (cor r * cis ϕ) = r2
    by (metis cis-inverse cis-neq-zero divide-complex-def cor-squared nonzero-mult-div-cancel-right power2-eq-square
semiring-normalization-rules(15))+
  ultimately
  have 1: poincare-line-cvec-cmat (of-complex-cvec (cor ry * cis ϑ)) (of-complex-cvec (cor r * cis ϕ)) = (?A, ?B,
cnj ?B, ?A)
    using den
    unfolding poincare-line-cvec-cmat-def of-complex-cvec-def Let-def prod.case
    by (simp add: field-simps)

  have 2: is-real ?A
    by simp
  let ?mix = cis ϑ * cis (-ϕ) - cis (-ϑ) * cis ϕ
  have is-imag ?mix
    using eq-minus-cnj-iff-imag[of ?mix]
    by simp
  hence Im ?mix ≠ 0
    using nondeg
    using complex.expand[of ?mix 0]
    by auto
  hence 3: Re ?A ≠ 0
    using ⟨r > 0⟩ ⟨ry > 0⟩
    by (simp add: sin-diff field-simps)

  have ?A ≠ 0
    using 2 3
    by auto
  hence 4: cor ?A ≠ 0
    using 2 3
    by (metis zero-complex.simps(1))

  have 5: ?ReB / ?A = (sin ϑ) / (2 * sin(ϕ - ϑ)) * (1/r + r) - (sin ϕ) / (2 * sin (ϕ - ϑ)) * (1/ry + ry)
    using ⟨ry > 0⟩ ⟨r > 0⟩
    apply (subst diff-divide-distrib)
    apply (subst add-frac-num, simp)
    apply (subst add-frac-num, simp)
    apply (simp add: power2-eq-square mult.commute)
    apply (simp add: field-simps)
  done

  have poincare-line-cvec-cmat (of-complex-cvec (cor ry * cis ϑ)) (of-complex-cvec (cor r * cis ϕ)) = (?A, ?B, cnj
?B, ?A) ∧

```

```

    is-real ?A ∧ Re ?A ≠ 0 ∧ ?A ≠ 0 ∧ cor ?A ≠ 0 ∧
    Re ?B = ?ReB ∧
    ?ReB / ?A = (sin ϑ) / (2 * sin(φ - ϑ)) * (1/r + r) - (sin φ) / (2 * sin(φ - ϑ)) * (1/ry + ry)
  using 1 2 3 4 5
  by auto
}
note ** = this

let ?Ayz = 2 * rz * ry * sin(φ - ϑ)
let ?Byz = i * (rz * cis φ * (1 + ry2) - ry * cis ϑ * (1 + rz2))
let ?ReByz = ry * (1 + rz2) * sin ϑ - rz * (1 + ry2) * sin φ
let ?Kz = (sin ϑ) / (2 * sin(φ - ϑ)) * (1/rz + rz) - (sin φ) / (2 * sin(φ - ϑ)) * (1/ry + ry)
have yz: poincare-line-cvec-cmat (of-complex-cvec (cor ry * cis ϑ)) (of-complex-cvec (cor rz * cis φ)) = (?Ayz,
?Byz, cnj ?Byz, ?Ayz)
  is-real ?Ayz Re ?Ayz ≠ 0 ?Ayz ≠ 0 cor ?Ayz ≠ 0 Re ?Byz = ?ReByz and Kz: ?ReByz / ?Ayz = ?Kz
  using **[OF <0 < rz>]
  by auto

let ?Ayu = 2 * ru * ry * sin(φ - ϑ)
let ?Byu = i * (ru * cis φ * (1 + ry2) - ry * cis ϑ * (1 + ru2))
let ?ReByu = ry * (1 + ru2) * sin ϑ - ru * (1 + ry2) * sin φ
let ?Ku = (sin ϑ) / (2 * sin(φ - ϑ)) * (1/ru + ru) - (sin φ) / (2 * sin(φ - ϑ)) * (1/ry + ry)
have yu: poincare-line-cvec-cmat (of-complex-cvec (cor ry * cis ϑ)) (of-complex-cvec (cor ru * cis φ)) = (?Ayu,
?Byu, cnj ?Byu, ?Ayu)
  is-real ?Ayu Re ?Ayu ≠ 0 ?Ayu ≠ 0 cor ?Ayu ≠ 0 Re ?Byu = ?ReByu and Ku: ?ReByu / ?Ayu = ?Ku
  using **[OF <0 < ru>]
  by auto

have ?Ayz ≠ 0
  using <sin(φ - ϑ) ≠ 0> <ry > 0> <rz > 0>
  by auto

have Re ?Byz / ?Ayz < -1
  using <intersects-x-axis-positive ?yz>
  * <y' = cor ry * cis ϑ> <z' = cor rz * cis φ> <u' = cor ru * cis φ>
  apply simp
  apply (transfer fixing: ry rz ru ϑ φ)
  apply (transfer fixing: ry rz ru ϑ φ)
proof-
  assume intersects-x-axis-positive-cmat (poincare-line-cvec-cmat (of-complex-cvec (cor ry * cis ϑ)) (of-complex-cvec
(cor rz * cis φ)))
  thus (ry * sin ϑ * (1 + rz2) - rz * sin φ * (1 + ry2)) / (2 * rz * ry * sin(φ - ϑ)) < - 1
  using yz
  by simp
qed

have ?ReByz / ?Ayz ≥ ?ReByu / ?Ayu
proof (cases sin φ > 0)
  case True
  hence sin ϑ < 0
  using <sin φ * sin ϑ < 0>
  by (smt mult-nonneg-nonneg)

have ?ReByz < 0
proof-
  have ry * (1 + rz2) * sin ϑ < 0
  using <ry > 0> <rz > 0>
  using <sin ϑ < 0>
  by (smt mult-pos-neg mult-pos-pos zero-less-power)
  moreover
  have rz * (1 + ry2) * sin φ > 0
  using <ry > 0> <rz > 0>
  using <sin φ > 0>
  by (smt mult-pos-neg mult-pos-pos zero-less-power)
  ultimately
  show ?thesis

```

```

    by simp
  qed
  have ?Ayz > 0
    using ⟨Re ?Byz / ?Ayz < -1⟩ ⟨Re ?Byz = ?ReByz⟩ ⟨?ReByz < 0⟩
    by (smt divide-less-0-iff)
  hence sin (φ - ϑ) > 0
    using ⟨ry > 0⟩ ⟨rz > 0⟩
    by (smt mult-pos-pos zero-less-mult-pos)

  have 1 / ru + ru ≥ 1 / rz + rz
    using Pasch-fun-mono[of ru rz] ⟨0 < ru⟩ ⟨ru ≤ rz⟩ ⟨rz < 1⟩
    by simp
  hence sin ϑ * (1 / ru + ru) ≤ sin ϑ * (1 / rz + rz)
    using ⟨sin ϑ < 0⟩
    by auto
  thus ?thesis
    using ⟨ru > 0⟩ ⟨rz > 0⟩ ⟨ru ≤ rz⟩ ⟨rz < 1⟩ ⟨?Ayz > 0⟩ ⟨sin (φ - ϑ) > 0⟩
    using divide-right-mono[of sin ϑ * (1 / ru + ru) sin ϑ * (1 / rz + rz) 2 * sin (φ - ϑ)]
    by (subst Kz, subst Ku) simp
next
  assume ¬ sin φ > 0
  hence sin φ < 0
    using ⟨sin φ * sin ϑ < 0⟩
    by (cases sin φ = 0, simp-all)
  hence sin ϑ > 0
    using ⟨sin φ * sin ϑ < 0⟩
    by (smt mult-nonpos-nonpos)
  have ?ReByz > 0
  proof-
    have ry * (1 + rz2) * sin ϑ > 0
      using ⟨ry > 0⟩ ⟨rz > 0⟩
      using ⟨sin ϑ > 0⟩
      by (smt mult-pos-neg mult-pos-pos zero-less-power)
    moreover
    have rz * (1 + ry2) * sin φ < 0
      using ⟨ry > 0⟩ ⟨rz > 0⟩
      using ⟨sin φ < 0⟩
      by (smt mult-pos-neg mult-pos-pos zero-less-power)
    ultimately
    show ?thesis
      by simp
  qed
  have ?Ayz < 0
    using ⟨Re ?Byz / ?Ayz < -1⟩ ⟨?Ayz ≠ 0⟩ ⟨Re ?Byz = ?ReByz⟩ ⟨?ReByz > 0⟩
    by (smt divide-less-0-iff)
  hence sin (φ - ϑ) < 0
    using ⟨ry > 0⟩ ⟨rz > 0⟩
    by (smt mult-nonneg-nonneg)

  have 1 / ru + ru ≥ 1 / rz + rz
    using Pasch-fun-mono[of ru rz] ⟨0 < ru⟩ ⟨ru ≤ rz⟩ ⟨rz < 1⟩
    by simp
  hence sin ϑ * (1 / ru + ru) ≥ sin ϑ * (1 / rz + rz)
    using ⟨sin ϑ > 0⟩
    by auto
  thus ?thesis
    using ⟨ru > 0⟩ ⟨rz > 0⟩ ⟨ru ≤ rz⟩ ⟨rz < 1⟩ ⟨?Ayz < 0⟩ ⟨sin (φ - ϑ) < 0⟩
    using divide-right-mono-neg[of sin ϑ * (1 / rz + rz) sin ϑ * (1 / ru + ru) 2 * sin (φ - ϑ)]
    by (subst Kz, subst Ku) simp
qed

  have intersects-x-axis-positive ?yu
    using * ⟨y' = cor ry * cis ϑ⟩ ⟨z' = cor rz * cis φ⟩ ⟨u' = cor ru * cis φ⟩
    apply simp
    apply (transfer fixing: ry rz ru ϑ φ)
    apply (transfer fixing: ry rz ru ϑ φ)

```

```

proof-
  have  $\text{Re } ?Byu / ?Ayu < -1$ 
    using  $\langle \text{Re } ?Byz / ?Ayz < -1 \rangle \langle ?ReByz / ?Ayz \geq ?ReByu / ?Ayu \rangle$ 
    by  $(\text{subst } (asm) \langle \text{Re } ?Byz = ?ReByz \rangle, \text{subst } \langle \text{Re } ?Byu = ?ReByu \rangle) \text{ simp}$ 
    thus  $\text{intersects-x-axis-positive-cmat } (\text{poincare-line-cvec-cmat } (\text{of-complex-cvec } (\text{cor } ry * cis \vartheta))) \text{ (of-complex-cvec } (cor\ ru * cis\ \varphi))$ 
    using  $yu$ 
    by  $\text{simp}$ 
qed

let  $?a = \text{calc-x-axis-intersection } ?yu$ 
have  $?a \in \text{positive-x-axis } ?a \in \text{circline-set } ?yu ?a \in \text{unit-disc}$ 
  using  $\langle \text{intersects-x-axis-positive } ?yu \rangle$ 
  using  $\text{intersects-x-axis-positive-iff}'[\text{of } ?yu] \langle y \neq u \rangle$ 
  by  $\text{auto}$ 

then obtain  $a'$  where  $a': ?a = \text{of-complex } a' \text{ is-real } a' \text{ Re } a' > 0 \text{ Re } a' < 1$ 
  unfolding  $\text{positive-x-axis-def circline-set-x-axis}$ 
  by  $(\text{auto simp add: cmod-eq-Re})$ 

have  $\text{intersects-x-axis } ?yz \text{ intersects-x-axis } ?yu$ 
  using  $\langle \text{intersects-x-axis-positive } ?yz \rangle \langle \text{intersects-x-axis-positive } ?yu \rangle$ 
  by  $\text{auto}$ 

show  $\exists a. a \in \text{unit-disc} \wedge \text{poincare-between } u\ a\ y \wedge \text{poincare-between } 0_h\ a \text{ (of-complex } v)$ 
proof  $(\text{rule-tac } x=?a \text{ in } exI, \text{ safe})$ 
  show  $\text{poincare-between } u\ ?a\ y$ 
    using  $\text{poincare-between-x-axis-intersection}[\text{of } y\ u\ ?a]$ 
    using  $\text{calc-x-axis-intersection}[OF \text{ is-poincare-line-poincare-line}[OF \langle y \neq u \rangle] \langle \text{intersects-x-axis } ?yu \rangle]$ 
    using  $\text{calc-x-axis-intersection-in-unit-disc}[OF \text{ is-poincare-line-poincare-line}[OF \langle y \neq u \rangle] \langle \text{intersects-x-axis } ?yu \rangle]$ 
    using  $\text{in-disc } \langle y \neq u \rangle \langle y \notin \text{circline-set } x\text{-axis} \rangle \langle u \notin \text{circline-set } x\text{-axis} \rangle$ 
    using  $* \langle \varphi = \text{Arg } u' \rangle \langle \vartheta = \text{Arg } y' \rangle \langle \varphi * \vartheta < 0 \rangle$ 
    by  $(\text{subst } \text{poincare-between-rev}, \text{auto simp add: mult.commute})$ 
  next
  show  $\text{poincare-between } 0_h\ ?a \text{ (of-complex } v)$ 
  proof-
    have  $-?ReByz / ?Ayz \leq -?ReByu / ?Ayu$ 
      using  $\langle ?ReByz / ?Ayz \geq ?ReByu / ?Ayu \rangle$ 
      by  $\text{linarith}$ 
    have  $\text{outward } ?yz\ ?yu$ 
      using  $* \langle y' = \text{cor } ry * cis\ \vartheta \rangle \langle z' = \text{cor } rz * cis\ \varphi \rangle \langle u' = \text{cor } ru * cis\ \varphi \rangle$ 
      apply  $\text{simp}$ 
      apply  $(\text{transfer fixing: } ry\ rz\ ru\ \vartheta\ \varphi)$ 
      apply  $(\text{transfer fixing: } ry\ rz\ ru\ \vartheta\ \varphi)$ 
      apply  $(\text{subst } yz\ yu) +$ 
      unfolding  $\text{outward-cmat-def}$ 
      apply  $(\text{simp only: Let-def prod.case})$ 
      apply  $(\text{subst } yz\ yu) +$ 
      using  $\langle -?ReByz / ?Ayz \leq -?ReByu / ?Ayu \rangle$ 
      by  $\text{simp}$ 
    hence  $\text{Re } a' \leq \text{Re } v$ 
      using  $\langle ?v = \text{of-complex } v \rangle$ 
      using  $\langle ?a = \text{of-complex } a' \rangle$ 
      using  $\langle \text{intersects-x-axis-positive } ?yz \rangle \langle \text{intersects-x-axis-positive } ?yu \rangle$ 
      using  $\text{outward}[OF \text{ is-poincare-line-poincare-line}[OF \langle y \neq z \rangle] \text{ is-poincare-line-poincare-line}[OF \langle y \neq u \rangle]]$ 
      by  $\text{simp}$ 
    thus  $?thesis$ 
      using  $\langle ?v = \text{of-complex } v \rangle$ 
      using  $\text{poincare-between-x-axis-0uv}[\text{of } \text{Re } a'\ \text{Re } v] a'\ v$ 
      by  $\text{simp}$ 
  qed
next
  show  $?a \in \text{unit-disc}$ 
  by  $\text{fact}$ 
qed
qed

```

```

next
  show  $x \in \text{unit-disc } v \in \text{unit-disc } x \neq v$ 
    using assms
    by auto
next
fix  $M x v$ 
let  $?Mx = \text{moebius-pt } M x$  and  $?Mv = \text{moebius-pt } M v$ 
assume 1:  $\text{unit-disc-fix } M x \in \text{unit-disc } v \in \text{unit-disc } x \neq v$ 
assume 2:  $?P ?Mx ?Mv$ 
show  $?P x v$ 
proof safe
  fix  $y z u$ 
  let  $?My = \text{moebius-pt } M y$  and  $?Mz = \text{moebius-pt } M z$  and  $?Mu = \text{moebius-pt } M u$ 
  assume  $\text{distinct } [x, y, z, u, v] \neg \text{poincare-collinear } \{x, y, z\} y \in \text{unit-disc } z \in \text{unit-disc } u \in \text{unit-disc}$ 
     $\text{poincare-between } x u z \text{poincare-between } y v z$ 
  hence  $\exists Ma. Ma \in \text{unit-disc} \wedge \text{poincare-between } ?Mu Ma ?My \wedge \text{poincare-between } ?Mx Ma ?Mv$ 
    using 1 2[rule-format, of ?My ?Mz ?Mu]
    by simp
  then obtain  $Ma$  where  $Ma: Ma \in \text{unit-disc} \text{poincare-between } ?Mu Ma ?My \wedge \text{poincare-between } ?Mx Ma ?Mv$ 
    by blast
  let  $?a = \text{moebius-pt } (-M) Ma$ 
  let  $?Ma = \text{moebius-pt } M ?a$ 
  have  $?Ma = Ma$ 
    by (metis moebius-pt-invert uminus-moebius-def)
  hence  $?Ma \in \text{unit-disc} \text{poincare-between } ?Mu ?Ma ?My \wedge \text{poincare-between } ?Mx ?Ma ?Mv$ 
    using  $Ma$ 
    by auto
  thus  $\exists a. a \in \text{unit-disc} \wedge \text{poincare-between } u a y \wedge \text{poincare-between } x a v$ 
    using  $\text{unit-disc-fix-moebius-inv}[OF \langle \text{unit-disc-fix } M \rangle \langle \text{unit-disc-fix } M \rangle \langle Ma \in \text{unit-disc} \rangle]$ 
    using  $\langle u \in \text{unit-disc} \rangle \langle v \in \text{unit-disc} \rangle \langle x \in \text{unit-disc} \rangle \langle y \in \text{unit-disc} \rangle$ 
    by (rule-tac x=?a in exI, simp del: moebius-pt-comp-inv-right)
qed
qed
thus ?thesis
  using assms
  by auto
qed

```

Pasch axiom, only degenerative cases.

lemma *Pasch-deg*:

```

assumes  $x \in \text{unit-disc}$  and  $y \in \text{unit-disc}$  and  $z \in \text{unit-disc}$  and  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
assumes  $\neg \text{distinct } [x, y, z, u, v] \vee \text{poincare-collinear } \{x, y, z\}$ 
assumes  $\text{poincare-between } x u z$  and  $\text{poincare-between } y v z$ 
shows  $\exists a. a \in \text{unit-disc} \wedge \text{poincare-between } u a y \wedge \text{poincare-between } x a v$ 
proof(cases poincare-collinear  $\{x, y, z\}$ )
  case True
  hence  $\text{poincare-between } x y z \vee \text{poincare-between } y x z \vee \text{poincare-between } y z x$ 
    using assms(1, 2, 3) poincare-collinear3-between poincare-between-rev by blast
  show ?thesis
  proof(cases poincare-between  $x y z$ )
    case True
    have  $\text{poincare-between } x y v$ 
      using True assms poincare-between-transitivity
      by (meson poincare-between-rev)
    thus ?thesis
      using assms(2)
      by (rule-tac x=y in exI, simp)
  next
  case False
  hence  $\text{poincare-between } y x z \vee \text{poincare-between } y z x$ 
    using  $\langle \text{poincare-between } x y z \vee \text{poincare-between } y x z \vee \text{poincare-between } y z x \rangle$ 
    by simp
  show ?thesis
  proof(cases poincare-between  $y x z$ )
    case True
    hence  $\text{poincare-between } u x y$ 

```



```

    using assms
  by (meson poincare-between-rev poincare-between-transitivity)
thus ?thesis
  using assms
  by (rule-tac  $x=x$  in exI, simp)
next
case False
hence poincare-between  $y z x$ 
  using  $\langle$ poincare-between  $y x z \vee$  poincare-between  $y z x$  $\rangle$ 
  by auto
hence poincare-between  $x z v$ 
  using assms
  by (meson poincare-between-rev poincare-between-transitivity)
hence poincare-between  $x u v$ 
  using assms poincare-between-transitivity poincare-between-rev
  by (smt poincare-between-sum-distances)
thus ?thesis
  using assms
  by (rule-tac  $x=u$  in exI, simp)
qed
qed
next
case False
hence  $\neg$  distinct  $[x, y, z, u, v]$ 
  using assms(6) by auto
show ?thesis
proof(cases  $u=z$ )
case True
thus ?thesis
  using assms
  apply(rule-tac  $x=v$  in exI)
  by(simp add:poincare-between-rev)
next
case False
hence  $x \neq z$ 
  using assms poincare-between-sandwich by blast
show ?thesis
proof(cases  $v=z$ )
case True
thus ?thesis
  using assms
  by (rule-tac  $x=u$  in exI, simp)
next
case False
hence  $y \neq z$ 
  using assms poincare-between-sandwich by blast
show ?thesis
proof(cases  $u = x$ )
case True
thus ?thesis
  using assms
  by (rule-tac  $x=x$  in exI, simp)
next
case False
have  $x \neq y$ 
  using assms  $\langle$  $\neg$  poincare-collinear  $\{x, y, z\}$  $\rangle$ 
  by fastforce
have  $x \neq v$ 
  using assms  $\langle$  $\neg$  poincare-collinear  $\{x, y, z\}$  $\rangle$ 
  by (metis insert-commute poincare-between-poincare-collinear)
have  $u \neq y$ 
  using assms  $\langle$  $\neg$  poincare-collinear  $\{x, y, z\}$  $\rangle$ 
  using poincare-between-poincare-collinear by blast
have  $u \neq v$ 
proof(rule ccontr)
  assume  $\neg u \neq v$ 

```

```

hence poincare-between  $x v z$ 
  using assms by auto
hence  $x \in \text{circline-set } (\text{poincare-line } z v)$ 
  using poincare-between-rev[ $of\ x\ v\ z$ ]
  using poincare-between-poincare-line-uvw[ $of\ z\ v\ x$ ]
  using assms  $\langle v \neq z \rangle$ 
  by auto
have  $y \in \text{circline-set } (\text{poincare-line } z v)$ 
  using assms  $\langle \neg u \neq v \rangle$ 
  using poincare-between-rev[ $of\ y\ v\ z$ ]
  using poincare-between-poincare-line-uvw[ $of\ z\ v\ y$ ]
  using assms  $\langle v \neq z \rangle$ 
  by auto
have  $z \in \text{circline-set } (\text{poincare-line } z v)$ 
  using ex-poincare-line-two-points[ $of\ z\ v$ ]  $\langle v \neq z \rangle$ 
  by auto
have is-poincare-line (poincare-line  $z v$ )
  using  $\langle v \neq z \rangle$ 
  by auto
hence poincare-collinear  $\{x, y, z\}$ 
  using  $\langle x \in \text{circline-set } (\text{poincare-line } z v) \rangle$ 
  using  $\langle y \in \text{circline-set } (\text{poincare-line } z v) \rangle$ 
  using  $\langle z \in \text{circline-set } (\text{poincare-line } z v) \rangle$ 
  unfolding poincare-collinear-def
  by (rule-tac  $x = \text{poincare-line } z v$  in exI, simp)
thus False
  using  $\langle \neg \text{poincare-collinear } \{x, y, z\} \rangle$  by simp
qed
have  $v = y$ 
  using  $\langle u \neq v \rangle \langle u \neq y \rangle \langle x \neq v \rangle \langle x \neq y \rangle \langle u \neq x \rangle \langle y \neq z \rangle \langle v \neq z \rangle \langle x \neq z \rangle \langle u \neq z \rangle$ 
  using  $\langle \neg \text{distinct } [x, y, z, u, v] \rangle$ 
  by auto
thus ?thesis
  using assms
  by (rule-tac  $x = y$  in exI, simp)
qed
qed
qed
qed

```

Axiom of Pasch

```

lemma Pasch:
  assumes  $x \in \text{unit-disc}$  and  $y \in \text{unit-disc}$  and  $z \in \text{unit-disc}$  and  $u \in \text{unit-disc}$  and  $v \in \text{unit-disc}$ 
  assumes poincare-between  $x u z$  and poincare-between  $y v z$ 
  shows  $\exists a. a \in \text{unit-disc} \wedge \text{poincare-between } u a y \wedge \text{poincare-between } x a v$ 
proof(cases distinct  $[x, y, z, u, v] \wedge \neg \text{poincare-collinear } \{x, y, z\}$ )
  case True
  thus ?thesis
    using assms Pasch-nondeg by auto
next
  case False
  thus ?thesis
    using assms Pasch-deg by auto
qed

```

11.2 Segment construction axiom

```

lemma segment-construction:
  assumes  $x \in \text{unit-disc}$  and  $y \in \text{unit-disc}$ 
  assumes  $a \in \text{unit-disc}$  and  $b \in \text{unit-disc}$ 
  shows  $\exists z. z \in \text{unit-disc} \wedge \text{poincare-between } x y z \wedge \text{poincare-distance } y z = \text{poincare-distance } a b$ 
proof-
  obtain  $d$  where  $d: d = \text{poincare-distance } a b$ 
  by auto
  have  $d \geq 0$ 
  using assms

```

```

by (simp add: d poincare-distance-ge0)

have  $\exists z. z \in \text{unit-disc} \wedge \text{poincare-between } x \ y \ z \wedge \text{poincare-distance } y \ z = d$  (is ?P x y)
proof (cases x = y)
case True
have  $\exists z. z \in \text{unit-disc} \wedge \text{poincare-distance } x \ z = d$ 
proof (rule wlog-zero)
show  $\exists z. z \in \text{unit-disc} \wedge \text{poincare-distance } 0_h \ z = d$ 
using ex-x-axis-poincare-distance-negative[of d]  $\langle d \geq 0 \rangle$ 
by blast
next
show  $x \in \text{unit-disc}$ 
by fact
next
fix a u
assume  $u \in \text{unit-disc} \text{ cmod } a < 1$ 
assume  $\exists z. z \in \text{unit-disc} \wedge \text{poincare-distance } (\text{moebius-pt } (\text{blaschke } a) \ u) \ z = d$ 
then obtain z where *:  $z \in \text{unit-disc} \text{ poincare-distance } (\text{moebius-pt } (\text{blaschke } a) \ u) \ z = d$ 
by auto
obtain z' where  $z': z = \text{moebius-pt } (\text{blaschke } a) \ z' \ z' \in \text{unit-disc}$ 
using  $\langle z \in \text{unit-disc} \rangle$ 
using unit-disc-fix-iff[of blaschke a]  $\langle \text{cmod } a < 1 \rangle$ 
using blaschke-unit-disc-fix[of a]
by blast

show  $\exists z. z \in \text{unit-disc} \wedge \text{poincare-distance } u \ z = d$ 
using * z'  $\langle u : \text{unit-disc} \rangle$ 
using blaschke-unit-disc-fix[of a]  $\langle \text{cmod } a < 1 \rangle$ 
by (rule-tac x=z' in exI, simp)
qed
thus ?thesis
using  $\langle x = y \rangle$ 
unfolding poincare-between-def
by auto
next
case False
show ?thesis
proof (rule wlog-positive-x-axis[where P= $\lambda y x. ?P \ x \ y$ ])
fix x
assume is-real x  $0 < \text{Re } x \ \text{Re } x < 1$ 

then obtain z where  $z: \text{is-real } z \ \text{Re } z \leq 0 - 1 < \text{Re } z \ \text{of-complex } z \in \text{unit-disc}$ 
of-complex z  $\in \text{unit-disc}$  of-complex z  $\in \text{circline-set } x\text{-axis} \ \text{poincare-distance } 0_h \ (\text{of-complex } z) = d$ 
using ex-x-axis-poincare-distance-negative[of d]  $\langle d \geq 0 \rangle$ 
by auto

have poincare-between (of-complex x)  $0_h$  (of-complex z)
proof (cases z = 0)
case True
thus ?thesis
unfolding poincare-between-def
by auto
next
case False
have  $x \neq 0$ 
using  $\langle \text{is-real } x \rangle \ \langle \text{Re } x > 0 \rangle$ 
by auto
thus ?thesis
using poincare-between-x-axis-u0v[of x z]
using z  $\langle \text{is-real } x \rangle \ \langle x \neq 0 \rangle \ \langle \text{Re } x > 0 \rangle$  False
using complex-eq-if-Re-eq mult-pos-neg
by fastforce
qed
thus ?P (of-complex x)  $0_h$ 
using  $\langle \text{poincare-distance } 0_h \ (\text{of-complex } z) = d \rangle \ \langle \text{of-complex } z \in \text{unit-disc} \rangle$ 
by blast

```

```

next
  show  $x \in \text{unit-disc } y \in \text{unit-disc}$ 
    by fact+
next
  show  $y \neq x$  using  $\langle x \neq y \rangle$  by simp
next
  fix  $M u v$ 
  assume  $\text{unit-disc-fix } M u \in \text{unit-disc } v \in \text{unit-disc } u \neq v$ 
  assume  $?P (\text{moebius-pt } M v) (\text{moebius-pt } M u)$ 
  then obtain  $z$  where  $*$ :  $z \in \text{unit-disc } \text{poincare-between } (\text{moebius-pt } M v) (\text{moebius-pt } M u) z \text{ poincare-distance}$ 
  ( $\text{moebius-pt } M u) z = d$ 
  by auto
  obtain  $z'$  where  $z': z = \text{moebius-pt } M z' z' \in \text{unit-disc}$ 
    using  $\langle z \in \text{unit-disc} \rangle$ 
    using  $\text{unit-disc-fix-iff}[of M] \langle \text{unit-disc-fix } M \rangle$ 
    by blast
  thus  $?P v u$ 
    using  $* \langle u \in \text{unit-disc} \rangle \langle v \in \text{unit-disc} \rangle \langle \text{unit-disc-fix } M \rangle$ 
    by auto
qed
qed
thus ?thesis
  using assms d
  by auto
qed

```

11.3 Five segment axiom

lemma *five-segment-axiom*:

```

assumes
  in-disc:  $x \in \text{unit-disc } y \in \text{unit-disc } z \in \text{unit-disc } u \in \text{unit-disc}$  and
  in-disc':  $x' \in \text{unit-disc } y' \in \text{unit-disc } z' \in \text{unit-disc } u' \in \text{unit-disc}$  and
   $x \neq y$  and
  betw:  $\text{poincare-between } x y z \text{ poincare-between } x' y' z'$  and
  xy:  $\text{poincare-distance } x y = \text{poincare-distance } x' y'$  and
  xu:  $\text{poincare-distance } x u = \text{poincare-distance } x' u'$  and
  yu:  $\text{poincare-distance } y u = \text{poincare-distance } y' u'$  and
  yz:  $\text{poincare-distance } y z = \text{poincare-distance } y' z'$ 
shows
   $\text{poincare-distance } z u = \text{poincare-distance } z' u'$ 
proof-
  from assms obtain  $M$  where
   $M: \text{unit-disc-fix-f } M M x = x' M u = u' M y = y'$ 
    using  $\text{unit-disc-fix-f-congruent-triangles}[of x y u]$ 
    by blast
  have  $M z = z'$ 
  proof (rule  $\text{unique-poincare-distance-on-ray}$ [where  $u=x'$  and  $v=y'$  and  $y=M z$  and  $z=z'$  and  $d=\text{poincare-distance } x z$ ])
    show  $0 \leq \text{poincare-distance } x z$ 
      using  $\text{poincare-distance-ge0 } \text{in-disc}$ 
      by simp
  next
    show  $x' \neq y'$ 
      using  $M \langle x \neq y \rangle$ 
      using  $\text{in-disc } \text{in-disc}' \text{poincare-distance-eq-0-iff } xy$ 
      by auto
  next
    show  $\text{poincare-distance } x' (M z) = \text{poincare-distance } x z$ 
      using  $M \text{in-disc}$ 
      unfolding  $\text{unit-disc-fix-f-def}$ 
      by auto
  next
    show  $M z \in \text{unit-disc}$ 
      using  $M \text{in-disc}$ 
      unfolding  $\text{unit-disc-fix-f-def}$ 
      by auto

```

```

next
  show poincare-distance  $x' z' = \textit{poincare-distance } x z$ 
    using xy yz betw
    using poincare-between-sum-distances[of x y z]
    using poincare-between-sum-distances[of x' y' z']
    using in-disc in-disc'
    by auto
next
  show poincare-between  $x' y' (M z)$ 
    using M
    using in-disc betw
    unfolding unit-disc-fix-f-def
    by auto
qed fact+
thus ?thesis
  using  $\langle \textit{unit-disc-fix-f } M \rangle$ 
  using in-disc in-disc'
   $\langle M u = u' \rangle$ 
  unfolding unit-disc-fix-f-def
  by auto
qed

```

11.4 Upper dimension axiom

```

lemma upper-dimension-axiom:
  assumes in-disc:  $x \in \textit{unit-disc } y \in \textit{unit-disc } z \in \textit{unit-disc } u \in \textit{unit-disc } v \in \textit{unit-disc}$ 
  assumes poincare-distance  $x u = \textit{poincare-distance } x v$ 
    poincare-distance  $y u = \textit{poincare-distance } y v$ 
    poincare-distance  $z u = \textit{poincare-distance } z v$ 
     $u \neq v$ 
  shows poincare-between  $x y z \vee \textit{poincare-between } y z x \vee \textit{poincare-between } z x y$ 
proof (cases  $x = y \vee y = z \vee x = z$ )
  case True
  thus ?thesis
    using in-disc
    by auto
next
  case False
  hence  $x \neq y \wedge x \neq z \wedge y \neq z$ 
    by auto
  let ?cong =  $\lambda a b a' b'. \textit{poincare-distance } a b = \textit{poincare-distance } a' b'$ 
  have  $\forall z u v. z \in \textit{unit-disc } \wedge u \in \textit{unit-disc } \wedge v \in \textit{unit-disc } \wedge$ 
    ?cong  $x u x v \wedge \textit{?cong } y u y v \wedge \textit{?cong } z u z v \wedge u \neq v \longrightarrow$ 
    poincare-collinear  $\{x, y, z\}$  (is ?P x y)
proof (rule wlog-positive-x-axis[where P=?P])
  fix x
  assume x: is-real  $x \ 0 < \textit{Re } x \ \textit{Re } x < 1$ 
  hence  $x \neq 0$ 
    by auto
  have  $0_h \in \textit{circline-set } x\text{-axis}$ 
    by simp
  show ?P  $0_h$  (of-complex x)
proof safe
  fix  $z u v$ 
  assume in-disc:  $z \in \textit{unit-disc } u \in \textit{unit-disc } v \in \textit{unit-disc}$ 
  then obtain  $z' u' v'$  where  $z = \textit{of-complex } z' \ u = \textit{of-complex } u' \ v = \textit{of-complex } v'$ 
    using inf-or-of-complex[of z] inf-or-of-complex[of u] inf-or-of-complex[of v]
    by auto

  assume cong: ?cong  $0_h u \ 0_h v \ \textit{?cong } (\textit{of-complex } x) u \ (\textit{of-complex } x) v \ \textit{?cong } z u z v \ u \neq v$ 

  let ?r0 = poincare-distance  $0_h u$  and
    ?rx = poincare-distance  $(\textit{of-complex } x) u$ 

  have ?r0  $> 0 \ \textit{?rx} > 0$ 
    using in-disc cong

```

```

using poincare-distance-eq-0-iff[of  $0_h$   $u$ ] poincare-distance-ge0[of  $0_h$   $u$ ]
using poincare-distance-eq-0-iff[of  $0_h$   $v$ ] poincare-distance-ge0[of  $0_h$   $v$ ]
using poincare-distance-eq-0-iff[of of-complex  $x$   $u$ ] poincare-distance-ge0[of of-complex  $x$   $u$ ]
using poincare-distance-eq-0-iff[of of-complex  $x$   $v$ ] poincare-distance-ge0[of of-complex  $x$   $v$ ]
using  $x$ 
by (auto simp add: cmod-eq-Re)

let  $?pc0 = \text{poincare-circle } 0_h \ ?r0$  and
 $?pcx = \text{poincare-circle (of-complex } x) \ ?rx$ 
have  $u \in ?pc0 \cap ?pcx \ v \in ?pc0 \cap ?pcx$ 
using in-disc cong
by (auto simp add: poincare-circle-def)
hence  $u = \text{conjugate } v$ 
using intersect-poincare-circles-x-axis[of  $0$   $x$   $?r0$   $?rx$   $u$   $v$ ]
using  $\langle x \neq 0 \rangle \langle u \neq v \rangle \langle ?r0 > 0 \rangle \langle ?rx > 0 \rangle$ 
by simp

let  $?ru = \text{poincare-distance } u \ z$ 
have  $?ru > 0$ 
using poincare-distance-ge0[of  $u$   $z$ ] in-disc
using cong
using poincare-distance-eq-0-iff[of  $z$   $u$ ] poincare-distance-eq-0-iff[of  $z$   $v$ ]
using poincare-distance-eq-0-iff
by force

have  $z \in \text{poincare-circle } u \ ?ru \cap \text{poincare-circle } v \ ?ru$ 
using cong in-disc
unfolding poincare-circle-def
by (simp add: poincare-distance-sym)

hence is-real  $z'$ 
using intersect-poincare-circles-conjugate-centers[of  $u$   $v$   $?ru$   $z$ ]  $\langle u = \text{conjugate } v \rangle \text{zuv}$ 
using in-disc  $\langle u \neq v \rangle \langle ?ru > 0 \rangle$ 
by simp

thus poincare-collinear  $\{0_h, \text{of-complex } x, z\}$ 
using poincare-line-0-real-is-x-axis[of of-complex  $x$ ]  $x \langle x \neq 0 \rangle \text{zuv} \langle 0_h \in \text{circline-set } x\text{-axis} \rangle$ 
unfolding poincare-collinear-def
by (rule-tac x=x-axis in exI, auto simp add: circline-set-x-axis)
qed
next
fix  $M \ x \ y$ 
assume 1: unit-disc-fix  $M \ x \in \text{unit-disc } y \in \text{unit-disc } x \neq y$ 
assume 2:  $?P (\text{moebius-pt } M \ x) (\text{moebius-pt } M \ y)$ 
show  $?P \ x \ y$ 
proof safe
fix  $z \ u \ v$ 
assume  $z \in \text{unit-disc } u \in \text{unit-disc } v \in \text{unit-disc}$ 
 $?cong \ x \ u \ x \ v \ ?cong \ y \ u \ y \ v \ ?cong \ z \ u \ z \ v \ u \neq v$ 
hence poincare-collinear  $\{\text{moebius-pt } M \ x, \text{moebius-pt } M \ y, \text{moebius-pt } M \ z\}$ 
using 1 2[rule-format, of moebius-pt M z moebius-pt M u moebius-pt M v]
by simp
then obtain  $p$  where is-poincare-line  $p \ \{\text{moebius-pt } M \ x, \text{moebius-pt } M \ y, \text{moebius-pt } M \ z\} \subseteq \text{circline-set } p$ 
unfolding poincare-collinear-def
by auto
thus poincare-collinear  $\{x, y, z\}$ 
using  $\langle \text{unit-disc-fix } M \rangle$ 
unfolding poincare-collinear-def
by (rule-tac x=moebius-circline (-M) p in exI, auto)
qed
qed fact+

thus ?thesis
using assms
using poincare-collinear3-between[of  $x \ y \ z$ ]
using poincare-between-rev

```

by *auto*
qed

11.5 Lower dimension axiom

lemma *lower-dimension-axiom*:

shows $\exists a \in \text{unit-disc. } \exists b \in \text{unit-disc. } \exists c \in \text{unit-disc.}$
 $\neg \text{poincare-between } a \ b \ c \wedge \neg \text{poincare-between } b \ c \ a \wedge \neg \text{poincare-between } c \ a \ b$

proof—

let $?u = \text{of-complex } (1/2)$ and $?v = \text{of-complex } (i/2)$
 have 1: $0_h \in \text{unit-disc}$ and 2: $?u \in \text{unit-disc}$ and 3: $?v \in \text{unit-disc}$
 by *simp-all*

have *: $\neg \text{poincare-collinear } \{0_h, ?u, ?v\}$

proof (rule *ccontr*)

assume $\neg ?thesis$

then obtain p where *is-poincare-line* $p \ \{0_h, ?u, ?v\} \subseteq \text{circline-set } p$
 unfolding *poincare-collinear-def*

by *auto*

moreover

have *of-complex* $(1 / 2) \neq \text{of-complex } (i / 2)$

using *of-complex-inj*

by *fastforce*

ultimately

have $0_h \in \text{circline-set } (\text{poincare-line } ?u \ ?v)$

using *unique-poincare-line*[*of ?u ?v p*]

by *auto*

thus *False*

unfolding *circline-set-def*

by *simp* (*transfer, transfer, simp add: vec-cnj-def*)

qed

show *?thesis*

apply (rule-tac $x=0_h$ in *beXI*, rule-tac $x=?u$ in *beXI*, rule-tac $x=?v$ in *beXI*)

apply (rule *ccontr, auto*)

using *

using *poincare-between-poincare-collinear*[*OF 1 2 3*]

using *poincare-between-poincare-collinear*[*OF 2 3 1*]

using *poincare-between-poincare-collinear*[*OF 3 1 2*]

by (*metis insert-commute*)+

qed

11.6 Negated Euclidean axiom

lemma *negated-euclidean-axiom-aux*:

assumes *on-circline* H (*of-complex* $(1/2 + i/2)$) and *is-poincare-line* H

assumes *intersects-x-axis-positive* H

shows $\neg \text{intersects-y-axis-positive } H$

using *assms*

proof (*transfer, transfer*)

fix H

assume *hh*: *hermitean* $H \wedge H \neq \text{mat-zero}$ *is-poincare-line-cmat* H

obtain $A \ B \ C \ D$ where $H = (A, B, C, D)$

by (*cases H, auto*)

hence *: *is-real* $A \ H = (A, B, \text{cnj } B, A) \ (\text{cmod } B)^2 > (\text{cmod } A)^2$

using *hermitean-elems*[*of A B C D*] *hh*

by *auto*

assume *intersects-x-axis-positive-cmat* H

hence $\text{Re } A \neq 0 \ \text{Re } B / \text{Re } A < -1$

using *

by *auto*

assume *on-circline-cmat-cvec* H (*of-complex-cvec* $(1 / 2 + i / 2)$)

hence $6*A + 4*Re \ B + 4*Im \ B = 0$

using *

unfolding *of-real-mult*

apply (*subst Re-express-cnj*[*of B*])

apply (*subst Im-express-cnj*[of B])
apply (*simp add: vec-cnj-def*)
apply (*simp add: field-simps*)
done
hence $6*A + 4*Re\ B + 4*Im\ B = 0$
by *simp*
hence $3*Re\ A + 2*Re\ B + 2*Im\ B = 0$
using $\langle is-real\ A \rangle$
by *simp*

hence $3/2 + Re\ B/Re\ A + Im\ B/Re\ A = 0$
using $\langle Re\ A \neq 0 \rangle$
by (*simp add: field-simps*)

hence $-Im\ B/Re\ A - 3/2 < -1$
using $\langle Re\ B / Re\ A < -1 \rangle$
by *simp*
hence $Im\ B/Re\ A > -1/2$
by (*simp add: field-simps*)
thus $\neg intersects-y-axis-positive-cmat\ H$
using $*$
by *simp*
qed

lemma *negated-euclidean-axiom:*

shows $\exists a\ b\ c\ d\ t.$

$a \in unit-disc \wedge b \in unit-disc \wedge c \in unit-disc \wedge d \in unit-disc \wedge t \in unit-disc \wedge$
poincare-between $a\ d\ t \wedge$ *poincare-between* $b\ d\ c \wedge a \neq d \wedge$

$(\forall x\ y. x \in unit-disc \wedge y \in unit-disc \wedge$

poincare-between $a\ b\ x \wedge$ *poincare-between* $x\ t\ y \longrightarrow \neg$ *poincare-between* $a\ c\ y)$

proof–

let $?a = 0_h$

let $?b = of-complex\ (1/2)$

let $?c = of-complex\ (i/2)$

let $?dl = (5 - sqrt\ 17) / 4$

let $?d = of-complex\ (?dl + i*?dl)$

let $?t = of-complex\ (1/2 + i/2)$

have $?dl \neq 0$

proof–

have $(sqrt\ 17)^2 \neq 5^2$

by *simp*

hence $sqrt\ 17 \neq 5$

by *force*

thus *?thesis*

by *simp*

qed

have $?d \neq ?a$

proof (*rule ccontr*)

assume $\neg ?thesis$

hence $?dl + i*?dl = 0$

by *simp*

hence $Re\ (?dl + i*?dl) = 0$

by *simp*

thus *False*

using $\langle ?dl \neq 0 \rangle$

by *simp*

qed

have $?dl > 0$

proof–

have $(sqrt\ 17)^2 < 5^2$

by (*simp add: power2-eq-square*)

hence $sqrt\ 17 < 5$

by (*rule power2-less-imp-less, simp*)


```

thus ?thesis
  by simp
qed

have ?a ≠ ?b
  by (metis divide-eq-0-iff of-complex-zero-iff zero-neq-numeral zero-neq-one)

have ?a ≠ ?c
  by (metis complex-i-not-zero divide-eq-0-iff of-complex-zero-iff zero-neq-numeral)

show ?thesis
proof (rule-tac x=?a in exI, rule-tac x=?b in exI, rule-tac x=?c in exI, rule-tac x=?d in exI, rule-tac x=?t in exI,
safe)

show ?a ∈ unit-disc ?b ∈ unit-disc ?c ∈ unit-disc ?t ∈ unit-disc
  by (auto simp add: cmod-def power2-eq-square)

have cmod-d: cmod (?dl + i*?dl) = ?dl * sqrt 2
  using ⟨?dl > 0⟩
  unfolding cmod-def
  by (simp add: real-sqrt-mult)

show ?d ∈ unit-disc
proof–
  have ?dl < 1 / sqrt 2
  proof–
    have 172 < (5 * sqrt 17)2
      by (simp add: field-simps)
    hence 17 < 5 * sqrt 17
      by (rule power2-less-imp-less, simp)
    hence ?dl2 < (1 / sqrt 2)2
      by (simp add: power2-eq-square field-simps)
    thus ?dl < 1 / sqrt 2
      by (rule power2-less-imp-less, simp)
  qed
thus ?thesis
  using cmod-d
  by (simp add: field-simps)
qed

have cmod-d: 1 - (cmod (to-complex ?d))2 = (-17 + 5*sqrt 17) / 4 (is - = ?cmod-d)
  apply (simp only: to-complex-of-complex)
  apply (subst cmod-d)
  apply (simp add: power-mult-distrib)
  apply (simp add: power2-eq-square field-simps)
  done

have cmod-d-c: (cmod (to-complex ?d - to-complex ?c))2 = (17 - 4*sqrt 17) / 4 (is - = ?cmod-dc)
  unfolding cmod-square
  by (simp add: field-simps)

have cmod-c: 1 - (cmod (to-complex ?c))2 = 3/4 (is - = ?cmod-c)
  by (simp add: power2-eq-square)

have xx: ∧ x::real. x + x = 2*x
  by simp

have cmod ((to-complex ?b) - (to-complex ?d)) = cmod ((to-complex ?d) - (to-complex ?c))
  by (simp add: cmod-def power2-eq-square field-simps)
moreover
have cmod (to-complex ?b) = cmod (to-complex ?c)
  by simp
ultimately
have *: poincare-distance-formula' (to-complex ?b) (to-complex ?d) =
  poincare-distance-formula' (to-complex ?d) (to-complex ?c)

```

```

unfolding poincare-distance-formula'-def
by simp

have **: poincare-distance-formula' (to-complex ?d) (to-complex ?c) = (sqrt 17) / 3
unfolding poincare-distance-formula'-def
proof (subst cmod-d, subst cmod-c, subst cmod-d-c)
  have (sqrt 17 * 15)2 ≠ 512
    by simp
  hence sqrt 17 * 15 ≠ 51
    by force
  hence sqrt 17 * 15 - 51 ≠ 0
    by simp

  have (5 * sqrt 17)2 ≠ 172
    by simp
  hence 5 * sqrt 17 ≠ 17
    by force
  hence ?cmod-d * ?cmod-c ≠ 0
    by simp
  hence 1 + 2 * (?cmod-dc / (?cmod-d * ?cmod-c)) = (?cmod-d * ?cmod-c + 2 * ?cmod-dc) / (?cmod-d * ?cmod-c)
    using add-frac-num[of ?cmod-d * ?cmod-c 2 * ?cmod-dc 1]
    by (simp add: field-simps)
  also have ... = (64 * (85 - sqrt 17 * 17)) / (64 * (sqrt 17 * 15 - 51))
    by (simp add: field-simps)
  also have ... = (85 - sqrt 17 * 17) / (sqrt 17 * 15 - 51)
    by (rule mult-divide-mult-cancel-left, simp)
  also have ... = sqrt 17 / 3
    by (subst frac-eq-eq, fact, simp, simp add: field-simps)
  finally
  show 1 + 2 * (?cmod-dc / (?cmod-d * ?cmod-c)) = sqrt 17 / 3
  .
qed

have sqrt 17 ≥ 3
proof-
  have (sqrt 17)2 ≥ 32
    by simp
  thus ?thesis
    by (rule power2-le-imp-le, simp)
qed

thus poincare-between ?b ?d ?c
  unfolding poincare-between-sum-distances[OF ‹?b ∈ unit-disc› ‹?d ∈ unit-disc› ‹?c ∈ unit-disc›]
  unfolding poincare-distance-formula[OF ‹?b ∈ unit-disc› ‹?d ∈ unit-disc›]
  unfolding poincare-distance-formula[OF ‹?d ∈ unit-disc› ‹?c ∈ unit-disc›]
  unfolding poincare-distance-formula[OF ‹?b ∈ unit-disc› ‹?c ∈ unit-disc›]
  unfolding poincare-distance-formula-def
  apply (subst *, subst xx, subst **, subst arcosh-double)
  apply (simp-all add: cmod-def power2-eq-square)
  done

show poincare-between ?a ?d ?t
proof (subst poincare-between-0uv[OF ‹?d ∈ unit-disc› ‹?t ∈ unit-disc› ‹?d ≠ ?a›])
  show ?t ≠ 0h
  proof (rule ccontr)
    assume ¬ ?thesis
    hence 1/2 + i/2 = 0
      by simp
    hence Re (1/2 + i/2) = 0
      by simp
    thus False
      by simp
  qed
next
  have 192 ≤ (5 * sqrt 17)2
    by simp
  hence 19 ≤ 5 * sqrt 17

```

```

  by (rule power2-le-imp-le, simp)
hence cmod (to-complex ?d) ≤ cmod (to-complex ?t)
  by (simp add: Let-def cmod-def power2-eq-square field-simps)
moreover
have Arg (to-complex ?d) = Arg (to-complex ?t)
proof-
  have 1: to-complex ?d = ((5 - sqrt 17) / 4) * (1 + i)
    by (simp add: field-simps)

  have 2: to-complex ?t = (cor (1/2)) * (1 + i)
    by (simp add: field-simps)

  have (sqrt 17)2 < 52
    by simp
  hence sqrt 17 < 5
    by (rule power2-less-imp-less, simp)
  hence 3: (5 - sqrt 17) / 4 > 0
    by simp

  have 4: (1::real) / 2 > 0
    by simp

  show ?thesis
    apply (subst 1, subst 2)
    apply (subst arg-mult-real-positive[OF 3])
    apply (subst arg-mult-real-positive[OF 4])
    by simp
qed
ultimately
show let d' = to-complex ?d; t' = to-complex ?t in Arg d' = Arg t' ∧ cmod d' ≤ cmod t'
  by simp
qed

show ?a = ?d ⇒ False
  using ⟨?d ≠ ?a⟩
  by simp

fix x y
assume x ∈ unit-disc y ∈ unit-disc

assume abx: poincare-between ?a ?b x
hence x ∈ circline-set x-axis
  using poincare-between-poincare-line-uvw[of ?a ?b x] ⟨x ∈ unit-disc⟩ ⟨?a ≠ ?b⟩
  using poincare-line-0-real-is-x-axis[of ?b]
  by (auto simp add: circline-set-x-axis)

have x ≠ 0h
  using abx poincare-between-sandwich[of ?a ?b] ⟨?a ≠ ?b⟩
  by auto

have x ∈ positive-x-axis
  using ⟨x ∈ circline-set x-axis⟩ ⟨x ≠ 0h⟩ ⟨x ∈ unit-disc⟩
  using abx poincare-between-x-axis-0uv[of 1/2 Re (to-complex x)]
  unfolding circline-set-x-axis positive-x-axis-def
  by (auto simp add: cmod-eq-Re abs-less-iff complex-eq-if-Re-eq)

assume acy: poincare-between ?a ?c y
hence y ∈ circline-set y-axis
  using poincare-between-poincare-line-uvw[of ?a ?c y] ⟨y ∈ unit-disc⟩ ⟨?a ≠ ?c⟩
  using poincare-line-0-imag-is-y-axis[of ?c]
  by (auto simp add: circline-set-y-axis)

have y ≠ 0h
  using acy poincare-between-sandwich[of ?a ?c] ⟨?a ≠ ?c⟩
  by auto

```

have $y \in \text{positive-}y\text{-axis}$
proof–
have $\bigwedge x. \llbracket x \neq 0; \text{poincare-between } 0_h \text{ (of-complex (i / 2)) (of-complex } x); \text{is-imag } x; -1 < \text{Im } x \rrbracket \implies 0 < \text{Im } x$
by (*smt add.left-neutral complex.expand divide-complex-def complex-eq divide-less-0-1-iff divide-less-eq-1-pos imaginary-unit.simps(1) mult.left-neutral of-real-1 of-real-add of-real-divide of-real-eq-0-iff one-add-one poincare-between-y-axis-0uv zero-complex.simps(1) zero-complex.simps(2) zero-less-divide-1-iff*)
thus *?thesis*
using $\langle y \in \text{circline-set } y\text{-axis} \rangle \langle y \neq 0_h \rangle \langle y \in \text{unit-disc} \rangle$
using *acy*
unfolding *circline-set-y-axis positive-y-axis-def*
by (*auto simp add: cmod-eq-Im abs-less-iff*)
qed

have $x \neq y$
using $\langle x \in \text{positive-}x\text{-axis} \rangle \langle y \in \text{positive-}y\text{-axis} \rangle$
unfolding *positive-x-axis-def positive-y-axis-def circline-set-x-axis circline-set-y-axis*
by *auto*

assume *xy: poincare-between x ?t y*

let *?xy = poincare-line x y*

have $?t \in \text{circline-set } ?xy$
using *xy poincare-between-poincare-line-uzv[OF $\langle x \neq y \rangle \langle x \in \text{unit-disc} \rangle \langle y \in \text{unit-disc} \rangle \langle ?t \in \text{unit-disc} \rangle$*
by *simp*

moreover

have $?xy \neq x\text{-axis}$
using *poincare-line-circline-set[OF $\langle x \neq y \rangle \langle y \in \text{positive-}y\text{-axis} \rangle$*
by (*auto simp add: circline-set-x-axis positive-y-axis-def*)
hence *intersects-x-axis-positive ?xy*
using *intersects-x-axis-positive-iff[of ?xy] $\langle x \neq y \rangle \langle x \in \text{unit-disc} \rangle \langle x \in \text{positive-}x\text{-axis} \rangle$*
by *auto*

moreover

have $?xy \neq y\text{-axis}$
using *poincare-line-circline-set[OF $\langle x \neq y \rangle \langle x \in \text{positive-}x\text{-axis} \rangle$*
by (*auto simp add: circline-set-y-axis positive-x-axis-def*)
hence *intersects-y-axis-positive ?xy*
using *intersects-y-axis-positive-iff[of ?xy] $\langle x \neq y \rangle \langle y \in \text{unit-disc} \rangle \langle y \in \text{positive-}y\text{-axis} \rangle$*
by *auto*

ultimately

show *False*
using *negated-euclidean-axiom-aux[of ?xy] $\langle x \neq y \rangle$*
unfolding *circline-set-def*
by *auto*

qed

qed

Alternate form of the Euclidean axiom – this one is much easier to prove

lemma *negated-euclidean-axiom'*:

shows $\exists a b c.$

$a \in \text{unit-disc} \wedge b \in \text{unit-disc} \wedge c \in \text{unit-disc} \wedge \neg(\text{poincare-collinear } \{a, b, c\}) \wedge$
 $\neg(\exists x. x \in \text{unit-disc} \wedge$
 $\text{poincare-distance } a x = \text{poincare-distance } b x \wedge$
 $\text{poincare-distance } a x = \text{poincare-distance } c x)$

proof–

let $?a = \text{of-complex (i/2)}$
let $?b = \text{of-complex (-i/2)}$
let $?c = \text{of-complex (1/5)}$

have $(i/2) \neq (-i/2)$

```

by simp
hence ?a ≠ ?b
  by (metis to-complex-of-complex)
have (i/2) ≠ (1/5)
  by simp
hence ?a ≠ ?c
  by (metis to-complex-of-complex)
have (-i/2) ≠ (1/5)
  by (simp add: minus-equation-iff)
hence ?b ≠ ?c
  by (metis to-complex-of-complex)

have ?a ∈ unit-disc ?b ∈ unit-disc ?c ∈ unit-disc
  by auto

moreover
have ¬(poincare-collinear {?a, ?b, ?c})
  unfolding poincare-collinear-def
proof(rule ccontr)
  assume ¬(∃ p. is-poincare-line p ∧ {?a, ?b, ?c} ⊆ circline-set p)
  then obtain p where is-poincare-line p ∧ {?a, ?b, ?c} ⊆ circline-set p
    by auto
  let ?ab = poincare-line ?a ?b
  have p = ?ab
    using ⟨is-poincare-line p ∧ {?a, ?b, ?c} ⊆ circline-set p⟩
    using unique-poincare-line[of ?a ?b] ⟨?a ≠ ?b⟩ ⟨?a ∈ unit-disc⟩ ⟨?b ∈ unit-disc⟩
    by auto
  have ?c ∉ circline-set ?ab
  proof(rule ccontr)
    assume ¬ ?c ∉ circline-set ?ab
    have poincare-between ?a 0h ?b
      unfolding poincare-between-def
      using cross-ratio-0inf by auto
    hence 0h ∈ circline-set ?ab
      using ⟨?a ≠ ?b⟩ ⟨?a ∈ unit-disc⟩ ⟨?b ∈ unit-disc⟩
      using poincare-between-poincare-line-uzv zero-in-unit-disc
      by blast
    hence ?ab = poincare-line 0h ?a
      using unique-poincare-line[of ?a ?b] ⟨?a ≠ ?b⟩ ⟨?a ∈ unit-disc⟩ ⟨?b ∈ unit-disc⟩
      using ⟨is-poincare-line p ∧ {?a, ?b, ?c} ⊆ circline-set p⟩
      using ⟨p = ?ab⟩ poincare-line-circline-set(1) unique-poincare-line
      by (metis add.inverse-neutral divide-minus-left of-complex-zero-iff zero-in-unit-disc)
    hence (i/2) * cnj(1/5) = cnj(i/2) * (1/5)
      using poincare-collinear-zero-iff[of (i/2) (1/5)]
      using ⟨?a ≠ ?c⟩ ⟨¬ ?c ∉ circline-set ?ab⟩ ⟨?a ∈ unit-disc⟩ ⟨?c ∈ unit-disc⟩ ⟨p = ?ab⟩
      using ⟨0h ∈ circline-set ?ab⟩ ⟨is-poincare-line p ∧ {?a, ?b, ?c} ⊆ circline-set p⟩
      using poincare-collinear-def by auto
    thus False
      by simp
  qed
  thus False
    using ⟨p = ?ab⟩ ⟨is-poincare-line p ∧ {?a, ?b, ?c} ⊆ circline-set p⟩
    by auto
qed

moreover

have ¬(∃ x. x ∈ unit-disc ∧
  poincare-distance ?a x = poincare-distance ?b x ∧
  poincare-distance ?a x = poincare-distance ?c x)
proof(rule ccontr)
  assume ¬ ?thesis
  then obtain x where x ∈ unit-disc poincare-distance ?a x = poincare-distance ?b x
    poincare-distance ?a x = poincare-distance ?c x
    by blast
  let ?x = to-complex x

```

```

have poincare-distance-formula' (i/2) ?x = poincare-distance-formula' (-i/2) ?x
  using ⟨poincare-distance ?a x = poincare-distance ?b x⟩
  using ⟨x ∈ unit-disc⟩ ⟨?a ∈ unit-disc⟩ ⟨?b ∈ unit-disc⟩
  by (metis cosh-dist to-complex-of-complex)
hence (cmod (i / 2 - ?x))2 = (cmod (- i / 2 - ?x))2
  unfolding poincare-distance-formula'-def
  apply (simp add:field-simps)
  using ⟨x ∈ unit-disc⟩ unit-disc-cmod-square-1 by fastforce
hence Im ?x = 0
  unfolding cmod-def
  by (simp add: power2-eq-iff)

have 1 - (Re ?x)2 ≠ 0
  using ⟨x ∈ unit-disc⟩ unit-disc-cmod-square-1
  using cmod-power2 by force
hence 24 - 24 * (Re ?x)2 ≠ 0
  by simp
have poincare-distance-formula' (i/2) ?x = poincare-distance-formula' (1/5) ?x
  using ⟨poincare-distance ?a x = poincare-distance ?c x⟩
  using ⟨x ∈ unit-disc⟩ ⟨?a ∈ unit-disc⟩ ⟨?c ∈ unit-disc⟩
  by (metis cosh-dist to-complex-of-complex)
hence (2 + 8 * (Re ?x)2) / (3 - 3 * (Re ?x)2) = 2 * (1 - Re ?x * 5)2 / (24 - 24 * (Re ?x)2) (is ?lhs = ?rhs)
  unfolding poincare-distance-formula'-def
  apply (simp add:field-simps)
  unfolding cmod-def
  using ⟨Im ?x = 0⟩
  by (simp add:field-simps)
hence *: ?lhs * (24 - 24 * (Re ?x)2) = ?rhs * (24 - 24 * (Re ?x)2)
  using ⟨(24 - 24 * (Re ?x)2) ≠ 0⟩
  by simp
have ?lhs * (24 - 24 * (Re ?x)2) = (2 + 8 * (Re ?x)2) * 8
  using ⟨(24 - 24 * (Re ?x)2) ≠ 0⟩ ⟨1 - (Re ?x)2 ≠ 0⟩
  by (simp add:field-simps)
have ?rhs * (24 - 24 * (Re ?x)2) = 2 * (1 - Re ?x * 5)2
  using ⟨(24 - 24 * (Re ?x)2) ≠ 0⟩ ⟨1 - (Re ?x)2 ≠ 0⟩
  by (simp add:field-simps)
hence (2 + 8 * (Re ?x)2) * 8 = 2 * (1 - Re ?x * 5)2
  using * ⟨?lhs * (24 - 24 * (Re ?x)2) = (2 + 8 * (Re ?x)2) * 8⟩
  by simp
hence 7 * (Re ?x)2 + 10 * (Re ?x) + 7 = 0
  by (simp add:field-simps comm-ring-1-class.power2-diff)
thus False
  using discriminant-iff[of 7 Re (to-complex x) 10 7] discrim-def[of 7 10 7]
  by auto
qed

```

```

ultimately show ?thesis
  apply (rule-tac x=?a in exI)
  apply (rule-tac x=?b in exI)
  apply (rule-tac x=?c in exI)
  by auto
qed

```

11.7 Continuity axiom

The set ϕ is on the left of the set ψ

abbreviation set-order where

set-order $A \varphi \psi \equiv \forall x \in \text{unit-disc}. \forall y \in \text{unit-disc}. \varphi x \wedge \psi y \longrightarrow \text{poincare-between } A x y$

The point B is between the sets ϕ and ψ

abbreviation point-between-sets where

point-between-sets $\varphi B \psi \equiv \forall x \in \text{unit-disc}. \forall y \in \text{unit-disc}. \varphi x \wedge \psi y \longrightarrow \text{poincare-between } x B y$

lemma continuity:

assumes $\exists A \in \text{unit-disc}. \text{set-order } A \varphi \psi$

```

shows  $\exists B \in \text{unit-disc. point-between-sets } \varphi B \psi$ 
proof (cases  $(\exists x0 \in \text{unit-disc. } \varphi x0) \wedge (\exists y0 \in \text{unit-disc. } \psi y0)$ )
  case False
  thus ?thesis
    using assms by blast
next
  case True
  then obtain Y0 where  $\psi Y0 Y0 \in \text{unit-disc}$ 
    by auto
  obtain A where  $*: A \in \text{unit-disc set-order } A \varphi \psi$ 
    using assms
    by auto
  show ?thesis
  proof(cases  $\forall x \in \text{unit-disc. } \varphi x \longrightarrow x = A$ )
    case True
    thus ?thesis
      using  $\langle A \in \text{unit-disc} \rangle$ 
      using poincare-between-nonstrict(1) by blast
  next
  case False
  then obtain X0 where  $\varphi X0 X0 \neq A X0 \in \text{unit-disc}$ 
    by auto
  have  $Y0 \neq A$ 
  proof(rule ccontr)
    assume  $\neg Y0 \neq A$ 
    hence  $\forall x \in \text{unit-disc. } \varphi x \longrightarrow \text{poincare-between } A x A$ 
      using  $* \langle \psi Y0 \rangle$ 
      by (cases A) force
    hence  $\forall x \in \text{unit-disc. } \varphi x \longrightarrow x = A$ 
      using  $* \text{poincare-between-sandwich}$  by blast
    thus False
      using False by auto
  qed

show ?thesis
proof (cases  $\exists B \in \text{unit-disc. } \varphi B \wedge \psi B$ )
  case True
  then obtain B where  $B \in \text{unit-disc } \varphi B \psi B$ 
    by auto
  hence  $\forall x \in \text{unit-disc. } \varphi x \longrightarrow \text{poincare-between } A x B$ 
    using  $*$  by auto
  have  $\forall y \in \text{unit-disc. } \psi y \longrightarrow \text{poincare-between } A B y$ 
    using  $* \langle B \in \text{unit-disc} \rangle \langle \varphi B \rangle$ 
    by auto

  show ?thesis
  proof(rule+)
    show  $B \in \text{unit-disc}$ 
      by fact
  next
  fix x y
  assume  $x \in \text{unit-disc } y \in \text{unit-disc } \varphi x \wedge \psi y$ 
  hence  $\text{poincare-between } A x B \text{poincare-between } A B y$ 
    using  $\langle \forall x \in \text{unit-disc. } \varphi x \longrightarrow \text{poincare-between } A x B \rangle$ 
    using  $\langle \forall y \in \text{unit-disc. } \psi y \longrightarrow \text{poincare-between } A B y \rangle$ 
    by simp+
  thus  $\text{poincare-between } x B y$ 
    using  $\langle x \in \text{unit-disc} \rangle \langle y \in \text{unit-disc} \rangle \langle B \in \text{unit-disc} \rangle \langle A \in \text{unit-disc} \rangle$ 
    using poincare-between-transitivity[of A x B y]
    by simp
  qed
next
  case False
  have  $\text{poincare-between } A X0 Y0$ 
    using  $\langle \varphi X0 \rangle \langle \psi Y0 \rangle * \langle Y0 \in \text{unit-disc} \rangle \langle X0 \in \text{unit-disc} \rangle$ 
    by auto

```

have $\forall \varphi. \forall \psi. \text{set-order } A \varphi \psi \wedge \neg (\exists B \in \text{unit-disc. } \varphi B \wedge \psi B) \wedge \varphi X0 \wedge$
 $(\exists y \in \text{unit-disc. } \psi y) \wedge (\exists x \in \text{unit-disc. } \varphi x)$
 $\longrightarrow (\exists B \in \text{unit-disc. } \text{point-between-sets } \varphi B \psi)$
(is ?P A X0)
proof (*rule wlog-positive-x-axis*[**where** $P=?P$])
show $A \in \text{unit-disc}$
by fact
next
show $X0 \in \text{unit-disc}$
by fact
next
show $A \neq X0$
using $\langle X0 \neq A \rangle$ **by simp**
next
fix $M u v$
let $?M = \lambda x. \text{moebius-pt } M x$
let $?Mu = ?M u$ **and** $?Mv = ?M v$
assume *hip*: $\text{unit-disc-fix } M u \in \text{unit-disc } v \in \text{unit-disc } u \neq v$
 $?P ?Mu ?Mv$
show $?P u v$
proof safe
fix $\varphi \psi x y$
assume $\text{set-order } u \varphi \psi \wedge \neg (\exists B \in \text{unit-disc. } \varphi B \wedge \psi B) \wedge \varphi v$
 $y \in \text{unit-disc } \psi y x \in \text{unit-disc } \varphi x$

let $?M\varphi = \lambda X'. \exists X. \varphi X \wedge ?M X = X'$
let $?M\psi = \lambda X'. \exists X. \psi X \wedge ?M X = X'$

obtain $M\varphi$ **where** $M\varphi = ?M\varphi$ **by simp**
obtain $M\psi$ **where** $M\psi = ?M\psi$ **by simp**

have $M\varphi ?Mv$
using $\langle \varphi v \rangle$ **using** $\langle M\varphi = ?M\varphi \rangle$
by blast
moreover
have $\neg (\exists B \in \text{unit-disc. } M\varphi B \wedge M\psi B)$
using $\langle \neg (\exists B \in \text{unit-disc. } \varphi B \wedge \psi B) \rangle$
using $\langle M\varphi = ?M\varphi \rangle \langle M\psi = ?M\psi \rangle$
by (*metis hip(1) moebius-pt-invert unit-disc-fix-discI unit-disc-fix-moebius-inv*)
moreover
have $\exists y \in \text{unit-disc. } M\psi y$
using $\langle y \in \text{unit-disc} \rangle \langle \psi y \rangle \langle M\psi = ?M\psi \rangle \langle \text{unit-disc-fix } M \rangle$
by auto
moreover
have $\text{set-order } ?Mu ?M\varphi ?M\psi$
proof (*(rule ballI)+, rule impI*)
fix $Mx My$
assume $Mx \in \text{unit-disc } My \in \text{unit-disc } ?M\varphi Mx \wedge ?M\psi My$
then obtain $x y$ **where** $\varphi x \wedge ?M x = Mx \wedge \psi y \wedge ?M y = My$
by blast

hence $x \in \text{unit-disc } y \in \text{unit-disc}$
using $\langle Mx \in \text{unit-disc} \rangle \langle My \in \text{unit-disc} \rangle \langle \text{unit-disc-fix } M \rangle$
by (*metis moebius-pt-comp-inv-left unit-disc-fix-discI unit-disc-fix-moebius-inv*)

hence *poincare-between* $u x y$
using $\langle \text{set-order } u \varphi \psi \rangle$
using $\langle Mx \in \text{unit-disc} \rangle \langle My \in \text{unit-disc} \rangle \langle \varphi x \wedge ?M x = Mx \rangle \langle \psi y \wedge ?M y = My \rangle$
by blast
then show *poincare-between* $?Mu Mx My$
using $\langle \varphi x \wedge ?M x = Mx \rangle \langle \psi y \wedge ?M y = My \rangle$
using $\langle x \in \text{unit-disc} \rangle \langle y \in \text{unit-disc} \rangle \langle u \in \text{unit-disc} \rangle \langle \text{unit-disc-fix } M \rangle$
using *unit-disc-fix-moebius-preserve-poincare-between* **by blast**
qed

hence $\text{set-order } ?Mu M\varphi M\psi$


```

    using ⟨Mφ = ?Mφ⟩ ⟨Mψ = ?Mψ⟩
    by simp
ultimately
have ∃ Mb ∈ unit-disc. point-between-sets Mφ Mb Mψ
    using hip(5)
    by blast
then obtain Mb where bbb:
    Mb ∈ unit-disc point-between-sets ?Mφ Mb ?Mψ
    using ⟨Mφ = ?Mφ⟩ ⟨Mψ = ?Mψ⟩
    by auto

let ?b = moebius-pt (moebius-inv M) Mb
show ∃ b ∈ unit-disc. point-between-sets φ b ψ
proof (rule-tac x=?b in beXI, (rule ballI)+, rule impI)
    fix x y
    assume x ∈ unit-disc y ∈ unit-disc φ x ∧ ψ y
    hence poincare-between u x y
        using ⟨set-order u φ ψ⟩
        by blast

let ?Mx = ?M x and ?My = ?M y

have ?Mφ ?Mx ?Mψ ?My
    using ⟨φ x ∧ ψ y⟩
    by blast+
have ?Mx ∈ unit-disc ?My ∈ unit-disc
    using ⟨x ∈ unit-disc⟩ ⟨unit-disc-fix M⟩ ⟨y ∈ unit-disc⟩
    by auto

hence poincare-between ?Mx Mb ?My
    using ⟨?Mφ ?Mx⟩ ⟨?Mψ ?My⟩ ⟨?Mx ∈ unit-disc⟩ ⟨?My ∈ unit-disc⟩ bbb
    by auto

then show poincare-between x ?b y
    using ⟨unit-disc-fix M⟩
    using ⟨x ∈ unit-disc⟩ ⟨y ∈ unit-disc⟩ ⟨Mb ∈ unit-disc⟩ ⟨?Mx ∈ unit-disc⟩ ⟨?My ∈ unit-disc⟩
    using unit-disc-fix-moebius-preserve-poincare-between[of M x ?b y]
    by auto
next
show ?b ∈ unit-disc
    using bbb ⟨unit-disc-fix M⟩
    by auto
qed
qed
next
fix X
assume xx: is-real X 0 < Re X Re X < 1
let ?X = of-complex X
show ?P 0h ?X
proof ((rule allI)+, rule impI, (erule conjE)+)
    fix φ ψ
    assume set-order 0h φ ψ ¬ (∃ B ∈ unit-disc. φ B ∧ ψ B) φ ?X
        ∃ y ∈ unit-disc. ψ y ∃ x ∈ unit-disc. φ x
    have ?X ∈ unit-disc
        using xx
        by (simp add: cmod-eq-Re)

have φpos: ∀ y ∈ unit-disc. ψ y ⟶ (is-real (to-complex y) ∧ Re (to-complex y) > 0)
proof (rule ballI, rule impI)
    fix y
    let ?y = to-complex y
    assume y ∈ unit-disc ψ y

hence poincare-between 0h ?X y
    using ⟨set-order 0h φ ψ⟩
    using ⟨?X ∈ unit-disc⟩ ⟨φ ?X⟩

```

by auto

thus $is\text{-}real\ ?y \wedge 0 < Re\ ?y$
using $\langle ?X \in unit\text{-}disc \rangle \langle y \in unit\text{-}disc \rangle$

by (*metis* (*mono-tags*, *opaque-lifting*) *arg-0-iff of-complex-zero-iff poincare-between-0uv poincare-between-sandwich to-complex-of-complex unit-disc-to-complex-inj zero-in-unit-disc*)

qed

have $\varphi_{noneg}: \forall x \in unit\text{-}disc. \varphi\ x \longrightarrow (is\text{-}real\ (to\text{-}complex\ x) \wedge Re\ (to\text{-}complex\ x) \geq 0)$

proof(*rule ballI*, *rule impI*)
fix x
assume $x \in unit\text{-}disc\ \varphi\ x$

obtain y **where** $y \in unit\text{-}disc\ \psi\ y$
using $\langle \exists y \in unit\text{-}disc. \psi\ y \rangle$ **by** *blast*

let $?x = to\text{-}complex\ x$ **and** $?y = to\text{-}complex\ y$

have $is\text{-}real\ ?y\ Re\ ?y > 0$
using $\psi_{pos}\ \langle \psi\ y \rangle \langle y \in unit\text{-}disc \rangle$
by auto

have *poincare-between* $0_h\ x\ y$
using $\langle set\text{-}order\ 0_h\ \varphi\ \psi \rangle$
using $\langle x \in unit\text{-}disc \rangle \langle \varphi\ x \rangle \langle y \in unit\text{-}disc \rangle \langle \psi\ y \rangle$
by auto

thus $is\text{-}real\ ?x \wedge 0 \leq Re\ ?x$
using $\langle x \in unit\text{-}disc \rangle \langle y \in unit\text{-}disc \rangle \langle is\text{-}real\ (to\text{-}complex\ y) \rangle \langle \psi\ y \rangle$
using $\langle set\text{-}order\ 0_h\ \varphi\ \psi \rangle$
using $\langle \varphi\ ?X \rangle \langle ?X \in unit\text{-}disc \rangle \langle Re\ ?y > 0 \rangle$
by (*metis* *arg-0-iff le-less of-complex-zero poincare-between-0uv to-complex-of-complex zero-complex.simps(1) zero-complex.simps(2)*)

qed

have $\varphi_{less\psi}: \forall x \in unit\text{-}disc. \forall y \in unit\text{-}disc. \varphi\ x \wedge \psi\ y \longrightarrow Re\ (to\text{-}complex\ x) < Re\ (to\text{-}complex\ y)$

proof(*(rule ballI)+*, *rule impI*)
fix $x\ y$
let $?x = to\text{-}complex\ x$ **and** $?y = to\text{-}complex\ y$
assume $x \in unit\text{-}disc\ y \in unit\text{-}disc\ \varphi\ x \wedge \psi\ y$

hence *poincare-between* $0_h\ x\ y$
using $\langle set\text{-}order\ 0_h\ \varphi\ \psi \rangle$
by auto

moreover

have $is\text{-}real\ ?x\ Re\ ?x \geq 0$
using φ_{noneg}
using $\langle x \in unit\text{-}disc \rangle \langle \varphi\ x \wedge \psi\ y \rangle$ **by auto**

moreover

have $is\text{-}real\ ?y\ Re\ ?y > 0$
using ψ_{pos}
using $\langle y \in unit\text{-}disc \rangle \langle \varphi\ x \wedge \psi\ y \rangle$ **by auto**

ultimately

have $Re\ ?x \leq Re\ ?y$
using $\langle x \in unit\text{-}disc \rangle \langle y \in unit\text{-}disc \rangle$
by (*metis* *Re-complex-of-real arg-0-iff le-less of-complex-zero poincare-between-0uv rcis-cmod-Arg rcis-zero-arg to-complex-of-complex*)

have $Re\ ?x \neq Re\ ?y$
using $\langle \varphi\ x \wedge \psi\ y \rangle \langle is\text{-}real\ ?x \rangle \langle is\text{-}real\ ?y \rangle$
using $\langle \neg (\exists B \in unit\text{-}disc. \varphi\ B \wedge \psi\ B) \rangle \langle x \in unit\text{-}disc \rangle \langle y \in unit\text{-}disc \rangle$
by (*metis* *complex.expand unit-disc-to-complex-inj*)

thus $Re\ ?x < Re\ ?y$
using $\langle Re\ ?x \leq Re\ ?y \rangle$ **by auto**

qed

have $\exists b \in \text{unit-disc}. \forall x \in \text{unit-disc}. \forall y \in \text{unit-disc}.$
 $\text{is-real } (\text{to-complex } b) \wedge$
 $(\varphi x \wedge \psi y \longrightarrow (\text{Re } (\text{to-complex } x) \leq \text{Re } (\text{to-complex } b) \wedge \text{Re } (\text{to-complex } b) \leq \text{Re } (\text{to-complex } y)))$

proof –
let $?Phi = \{x. (\text{of-complex } (\text{cor } x)) \in \text{unit-disc} \wedge \varphi (\text{of-complex } (\text{cor } x))\}$

have $\forall x \in \text{unit-disc}. \varphi x \longrightarrow \text{Re } (\text{to-complex } x) \leq \text{Sup } ?Phi$

proof(*safe*)
fix x
let $?x = \text{to-complex } x$
assume $x \in \text{unit-disc} \varphi x$
hence $\text{is-real } ?x \text{ Re } ?x \geq 0$
using φnoneg
by *auto*
hence $\text{cor } (\text{Re } ?x) = ?x$
using *complex-of-real-Re* **by** *blast*
hence $\text{of-complex } (\text{cor } (\text{Re } ?x)) \in \text{unit-disc}$
using $\langle x \in \text{unit-disc} \rangle$
by (*metis inf-notin-unit-disc of-complex-to-complex*)
moreover
have $\varphi (\text{of-complex } (\text{cor } (\text{Re } ?x)))$
using $\langle \text{cor } (\text{Re } ?x) = ?x \rangle \langle \varphi x \rangle \langle x \in \text{unit-disc} \rangle$
by (*metis inf-notin-unit-disc of-complex-to-complex*)
ultimately
have $\text{Re } ?x \in ?Phi$
by *auto*

have $\exists M. \forall x \in ?Phi. x \leq M$
using $\varphi\text{less}\psi$
using $\langle \exists y \in \text{unit-disc}. \psi y \rangle$
by (*metis (mono-tags, lifting) Re-complex-of-real le-less mem-Collect-eq to-complex-of-complex*)

thus $\text{Re } ?x \leq \text{Sup } ?Phi$
using *cSup-upper*[$\text{of } \text{Re } ?x ?Phi$]
unfolding *bdd-above-def*
using $\langle \text{Re } ?x \in ?Phi \rangle$
by *auto*

qed

have $\forall y \in \text{unit-disc}. \psi y \longrightarrow \text{Sup } ?Phi \leq \text{Re } (\text{to-complex } y)$

proof (*safe*)
fix y
let $?y = \text{to-complex } y$
assume $\psi y y \in \text{unit-disc}$
show $\text{Sup } ?Phi \leq \text{Re } ?y$
proof (*rule ccontr*)
assume $\neg ?thesis$
hence $\text{Re } ?y < \text{Sup } ?Phi$
by *auto*

have $\exists x. \varphi (\text{of-complex } (\text{cor } x)) \wedge (\text{of-complex } (\text{cor } x)) \in \text{unit-disc}$

proof –
obtain x' **where** $x' \in \text{unit-disc} \varphi x'$
using $\langle \exists x \in \text{unit-disc}. \varphi x \rangle$ **by** *blast*
let $?x' = \text{to-complex } x'$
have $\text{is-real } ?x'$
using $\langle x' \in \text{unit-disc} \rangle \langle \varphi x' \rangle$
using φnoneg
by *auto*
hence $\text{cor } (\text{Re } ?x') = ?x'$
using *complex-of-real-Re* **by** *blast*
hence $x' = \text{of-complex } (\text{cor } (\text{Re } ?x'))$
using $\langle x' \in \text{unit-disc} \rangle$
by (*metis inf-notin-unit-disc of-complex-to-complex*)
show $?thesis$

apply (*rule-tac* $x=Re \ ?x'$ **in** exI)
using $\langle x' \in unit-disc \rangle$
apply (*subst* (*asm*) $\langle x' = of-complex (cor (Re \ ?x')) \rangle$, *simp*)
using $\langle \varphi \ x' \rangle$
by (*subst* (*asm*) (2) $\langle x' = of-complex (cor (Re \ ?x')) \rangle$, *simp*)
qed

hence $?Phi \neq \{\}$
by *auto*

then obtain x **where** $\varphi (of-complex (cor \ x)) \ Re \ ?y < x$
 $(of-complex (cor \ x)) \in unit-disc$

using $\langle Re \ ?y < Sup \ ?Phi \rangle$
using *less-cSupE*[*of* $Re \ ?y \ ?Phi$]
by *auto*

moreover

have $Re \ ?y < Re (to-complex (of-complex (cor \ x)))$

using $\langle Re \ ?y < x \rangle$

by *simp*

ultimately

show *False*

using $\langle \varphi less\psi \rangle$

using $\langle \psi \ y \rangle \langle y \in unit-disc \rangle$

by (*metis* *less-not-sym*)

qed

qed

thus *?thesis*

using $\langle \forall x \in unit-disc. \varphi \ x \longrightarrow Re (to-complex \ x) \leq Sup \ ?Phi \rangle$

apply (*rule-tac* $x=(of-complex (cor (Sup \ ?Phi)))$ **in** bxI , *simp*)

using $\langle \exists y \in unit-disc. \psi \ y \rangle \langle \varphi \ ?X \rangle \langle ?X \in unit-disc \rangle$

using $\langle \forall y \in unit-disc. \psi \ y \longrightarrow is-real (to-complex \ y) \wedge 0 < Re (to-complex \ y) \rangle$

by (*smt* *complex-of-real-Re* *inf-notin-unit-disc* *norm-of-real* *of-complex-to-complex* *to-complex-of-complex* *unit-disc-iff-cmod-lt-1* $xx(2)$)

qed

then obtain B **where** $B \in unit-disc$ *is-real* (*to-complex* B)

$\forall x \in unit-disc. \forall y \in unit-disc. \varphi \ x \wedge \psi \ y \longrightarrow Re (to-complex \ x) \leq Re (to-complex \ B) \wedge$

$Re (to-complex \ B) \leq Re (to-complex \ y)$

by *blast*

show $\exists b \in unit-disc.$ *point-between-sets* $\varphi \ b \ \psi$

proof (*rule-tac* $x=B$ **in** bxI)

show $B \in unit-disc$

by *fact*

next

show *point-between-sets* $\varphi \ B \ \psi$

proof ((*rule* *ballI*)₊, *rule* *impI*)

fix $x \ y$

let $?x = to-complex \ x$ **and** $?y = to-complex \ y$ **and** $?B = to-complex \ B$

assume $x \in unit-disc \ y \in unit-disc \ \varphi \ x \wedge \psi \ y$

hence $Re \ ?x \leq Re \ ?B \wedge Re \ ?B \leq Re \ ?y$

using $\langle \forall x \in unit-disc. \forall y \in unit-disc. \varphi \ x \wedge \psi \ y \longrightarrow Re (to-complex \ x) \leq Re \ ?B \wedge$
 $Re (to-complex \ B) \leq Re (to-complex \ y) \rangle$

by *auto*

moreover

have *is-real* $?x \ Re \ ?x \geq 0$

using *nonneg*

using $\langle x \in unit-disc \rangle \langle \varphi \ x \wedge \psi \ y \rangle$

by *auto*

moreover

have *is-real* $?y \ Re \ ?y > 0$

using *ψpos*

using $\langle y \in unit-disc \rangle \langle \varphi \ x \wedge \psi \ y \rangle$

by *auto*

```

moreover
have  $\text{cor } (\text{Re } ?x) = ?x$ 
  using  $\text{complex-of-real-Re } \langle \text{is-real } ?x \rangle$  by blast
hence  $x = \text{of-complex } (\text{cor } (\text{Re } ?x))$ 
  using  $\langle x \in \text{unit-disc} \rangle$ 
  by  $(\text{metis inf-notin-unit-disc of-complex-to-complex})$ 
moreover
have  $\text{cor } (\text{Re } ?y) = ?y$ 
  using  $\text{complex-of-real-Re } \langle \text{is-real } ?y \rangle$  by blast
hence  $y = \text{of-complex } (\text{cor } (\text{Re } ?y))$ 
  using  $\langle y \in \text{unit-disc} \rangle$ 
  by  $(\text{metis inf-notin-unit-disc of-complex-to-complex})$ 
moreover
have  $\text{cor } (\text{Re } ?B) = ?B$ 
  using  $\text{complex-of-real-Re } \langle \text{is-real } (\text{to-complex } B) \rangle$  by blast
hence  $B = \text{of-complex } (\text{cor } (\text{Re } ?B))$ 
  using  $\langle B \in \text{unit-disc} \rangle$ 
  by  $(\text{metis inf-notin-unit-disc of-complex-to-complex})$ 
ultimately
show poincare-between  $x B y$ 
  using  $\langle \text{is-real } (\text{to-complex } B) \rangle \langle x \in \text{unit-disc} \rangle \langle y \in \text{unit-disc} \rangle \langle B \in \text{unit-disc} \rangle$ 
  using poincare-between-x-axis-uvw $[\text{of Re } (\text{to-complex } x) \text{ Re } (\text{to-complex } B) \text{ Re } (\text{to-complex } y)]$ 
by  $(\text{smt Re-complex-of-real arg-0-iff poincare-between-nonstrict}(1) \text{rcis-cmod-Arg rcis-zero-arg unit-disc-iff-cmod-lt-1})$ 
qed
qed
qed
qed
thus ?thesis
  using  $\text{False } \langle \varphi X0 \rangle \langle \psi Y0 \rangle * \langle Y0 \in \text{unit-disc} \rangle \langle X0 \in \text{unit-disc} \rangle$ 
  by auto
qed
qed
qed

```

11.8 Limiting parallels axiom

Auxiliary definitions

definition *poincare-on-line* **where**

$\text{poincare-on-line } p a b \longleftrightarrow \text{poincare-collinear } \{p, a, b\}$

definition *poincare-on-ray* **where**

$\text{poincare-on-ray } p a b \longleftrightarrow \text{poincare-between } a p b \vee \text{poincare-between } a b p$

definition *poincare-in-angle* **where**

$\text{poincare-in-angle } p a b c \longleftrightarrow$

$b \neq a \wedge b \neq c \wedge p \neq b \wedge (\exists x \in \text{unit-disc. } \text{poincare-between } a x c \wedge x \neq a \wedge x \neq c \wedge \text{poincare-on-ray } p b x)$

definition *poincare-ray-meets-line* **where**

$\text{poincare-ray-meets-line } a b c d \longleftrightarrow (\exists x \in \text{unit-disc. } \text{poincare-on-ray } x a b \wedge \text{poincare-on-line } x c d)$

All points on ray are collinear

lemma *poincare-on-ray-poincare-collinear*:

assumes $p \in \text{unit-disc}$ **and** $a \in \text{unit-disc}$ **and** $b \in \text{unit-disc}$ **and** $\text{poincare-on-ray } p a b$

shows $\text{poincare-collinear } \{p, a, b\}$

using *assms poincare-between-poincare-collinear*

unfolding *poincare-on-ray-def*

by $(\text{metis insert-commute})$

H-isometries preserve all defined auxiliary relations

lemma *unit-disc-fix-preserves-poincare-on-line* [*simp*]:

assumes $\text{unit-disc-fix } M$ **and** $p \in \text{unit-disc}$ $a \in \text{unit-disc}$ $b \in \text{unit-disc}$

shows $\text{poincare-on-line } (\text{moebius-pt } M p) (\text{moebius-pt } M a) (\text{moebius-pt } M b) \longleftrightarrow \text{poincare-on-line } p a b$

using *assms*

unfolding *poincare-on-line-def*

by *auto*

lemma *unit-disc-fix-preserves-poincare-on-ray* [simp]:
assumes *unit-disc-fix* $M p \in \text{unit-disc } a \in \text{unit-disc } b \in \text{unit-disc}$
shows *poincare-on-ray* (moebius-pt $M p$) (moebius-pt $M a$) (moebius-pt $M b$) \longleftrightarrow *poincare-on-ray* $p a b$
using *assms*
unfolding *poincare-on-ray-def*
by *auto*

lemma *unit-disc-fix-preserves-poincare-in-angle* [simp]:
assumes *unit-disc-fix* $M p \in \text{unit-disc } a \in \text{unit-disc } b \in \text{unit-disc } c \in \text{unit-disc}$
shows *poincare-in-angle* (moebius-pt $M p$) (moebius-pt $M a$) (moebius-pt $M b$) (moebius-pt $M c$) \longleftrightarrow *poincare-in-angle* $p a b c$ (is ?lhs \longleftrightarrow ?rhs)
proof
assume ?lhs
then obtain Mx **where** *: $Mx \in \text{unit-disc}$
poincare-between (moebius-pt $M a$) Mx (moebius-pt $M c$)
 $Mx \neq \text{moebius-pt } M a$ $Mx \neq \text{moebius-pt } M c$ *poincare-on-ray* (moebius-pt $M p$) (moebius-pt $M b$) Mx
 $\text{moebius-pt } M b \neq \text{moebius-pt } M a$ $\text{moebius-pt } M b \neq \text{moebius-pt } M c$ $\text{moebius-pt } M p \neq \text{moebius-pt } M b$
unfolding *poincare-in-angle-def*
by *auto*
obtain x **where** $Mx = \text{moebius-pt } M x$ $x \in \text{unit-disc}$
by (*metis* *(1) *assms*(1) *image-iff unit-disc-fix-iff*)
thus ?rhs
using * *assms*
unfolding *poincare-in-angle-def*
by *auto*

next
assume ?rhs
then obtain x **where** *: $x \in \text{unit-disc}$
poincare-between $a x c$
 $x \neq a$ $x \neq c$ *poincare-on-ray* $p b x$
 $b \neq a$ $b \neq c$ $p \neq b$
unfolding *poincare-in-angle-def*
by *auto*
thus ?lhs
using *assms*
unfolding *poincare-in-angle-def*
by *auto* (*rule-tac* $x = \text{moebius-pt } M x$ **in** *be* xI , *auto*)

qed

lemma *unit-disc-fix-preserves-poincare-ray-meets-line* [simp]:
assumes *unit-disc-fix* $M a \in \text{unit-disc } b \in \text{unit-disc } c \in \text{unit-disc } d \in \text{unit-disc}$
shows *poincare-ray-meets-line* (moebius-pt $M a$) (moebius-pt $M b$) (moebius-pt $M c$) (moebius-pt $M d$) \longleftrightarrow *poincare-ray-meets-line* $a b c d$ (is ?lhs \longleftrightarrow ?rhs)
proof
assume ?lhs
then obtain Mx **where** *: $Mx \in \text{unit-disc}$ *poincare-on-ray* Mx (moebius-pt $M a$) (moebius-pt $M b$)
poincare-on-line Mx (moebius-pt $M c$) (moebius-pt $M d$)
unfolding *poincare-ray-meets-line-def*
by *auto*
obtain x **where** $Mx = \text{moebius-pt } M x$ $x \in \text{unit-disc}$
by (*metis* *(1) *assms*(1) *image-iff unit-disc-fix-iff*)
thus ?rhs
using *assms* *
unfolding *poincare-ray-meets-line-def* *poincare-on-line-def*
by *auto*

next
assume ?rhs
then obtain x **where** *: $x \in \text{unit-disc}$ *poincare-on-ray* $x a b$
poincare-on-line $x c d$
unfolding *poincare-ray-meets-line-def*
by *auto*
thus ?lhs
using *assms* *
unfolding *poincare-ray-meets-line-def* *poincare-on-line-def*
by *auto* (*rule-tac* $x = \text{moebius-pt } M x$ **in** *be* xI , *auto*)

qed

H-lines that intersect on the absolute do not meet (they do not share a common h-point)

lemma *tangent-not-meet*:

assumes $x1 \in \text{unit-disc}$ **and** $x2 \in \text{unit-disc}$ **and** $x1 \neq x2$ **and** $\neg \text{poincare-collinear } \{0_h, x1, x2\}$

assumes $i \in \text{ideal-points } (\text{poincare-line } x1 \ x2)$ $a \in \text{unit-disc}$ $a \neq 0_h$ $\text{poincare-collinear } \{0_h, a, i\}$

shows $\neg \text{poincare-ray-meets-line } 0_h \ a \ x1 \ x2$

proof (rule *ccontr*)

assume $\neg ?thesis$

then obtain x **where** $x \in \text{unit-disc}$ $\text{poincare-on-ray } x \ 0_h \ a$ $\text{poincare-collinear } \{x, x1, x2\}$

unfolding *poincare-ray-meets-line-def* *poincare-on-line-def*

by *auto*

have $\text{poincare-collinear } \{0_h, a, x\}$

using $\langle \text{poincare-on-ray } x \ 0_h \ a \rangle \langle x \in \text{unit-disc} \rangle \langle a \in \text{unit-disc} \rangle$

by (*meson* *poincare-between-poincare-collinear* *poincare-between-rev* *poincare-on-ray-def* *poincare-on-ray-poincare-collinear* *zero-in-unit-disc*)

have $x \neq 0_h$

using $\langle \neg \text{poincare-collinear } \{0_h, x1, x2\} \rangle \langle \text{poincare-collinear } \{x, x1, x2\} \rangle$

unfolding *poincare-collinear-def*

by (*auto simp add: assms(2) assms(3) poincare-between-rev*)

let $?l1 = \text{poincare-line } 0_h \ a$

let $?l2 = \text{poincare-line } x1 \ x2$

have $i \in \text{circline-set unit-circle}$

using $\langle i \in \text{ideal-points } (\text{poincare-line } x1 \ x2) \rangle$

using *assms(3) ideal-points-on-unit-circle is-poincare-line-poincare-line* **by** *blast*

have $i \in \text{circline-set } ?l1$

using $\langle \text{poincare-collinear } \{0_h, a, i\} \rangle$

unfolding *poincare-collinear-def*

using $\langle a \in \text{unit-disc} \rangle \langle a \neq 0_h \rangle$

by (*metis insert-subset unique-poincare-line zero-in-unit-disc*)

moreover

have $x \in \text{circline-set } ?l1$

using $\langle a \in \text{unit-disc} \rangle \langle a \neq 0_h \rangle \langle \text{poincare-collinear } \{0_h, a, x\} \rangle \langle x \in \text{unit-disc} \rangle$

by (*metis* *poincare-collinear3-between* *poincare-between-poincare-line-uvw* *poincare-between-poincare-line-uvw* *poincare-line-sym* *zero-in-unit-disc*)

moreover

have *inversion* $x \in \text{circline-set } ?l1$

using $\langle \text{poincare-collinear } \{0_h, a, x\} \rangle$

using *poincare-line-inversion-full*[of $0_h \ a \ x$] $\langle a \in \text{unit-disc} \rangle \langle a \neq 0_h \rangle \langle x \in \text{unit-disc} \rangle$

by (*metis* *poincare-collinear3-between* *is-poincare-line-inverse-point* *is-poincare-line-poincare-line* *poincare-between-poincare-line-uvw* *poincare-between-poincare-line-uvw* *poincare-line-sym* *zero-in-unit-disc*)

moreover

have $x \in \text{circline-set } ?l2$

using $\langle \text{poincare-collinear } \{x, x1, x2\} \rangle \langle x1 \neq x2 \rangle \langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle \langle x \in \text{unit-disc} \rangle$

by (*metis* *insert-commute inversion-noteq-unit-disc* *poincare-between-poincare-line-uvw* *poincare-between-poincare-line-uvw* *poincare-collinear3-iff* *poincare-line-sym-general*)

moreover

hence *inversion* $x \in \text{circline-set } ?l2$

using $\langle x1 \neq x2 \rangle \langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle \langle x \in \text{unit-disc} \rangle$

using *poincare-line-inversion-full*[of $x1 \ x2 \ x$]

unfolding *circline-set-def*

by *auto*

moreover

have $i \in \text{circline-set } ?l2$
using $\langle x1 \neq x2 \rangle \langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle$
using $\langle i \in \text{ideal-points } ?l2 \rangle$
by (simp add: ideal-points-on-circline)

moreover

have $x \neq \text{inversion } x$
using $\langle x \in \text{unit-disc} \rangle$
using inversion-noteq-unit-disc by fastforce

moreover

have $x \neq i$
using $\langle x \in \text{unit-disc} \rangle$
using $\langle i \in \text{circline-set unit-circle} \rangle \text{circline-set-def inversion-noteq-unit-disc}$
by fastforce+

moreover

have inversion $x \neq i$
using $\langle i \in \text{circline-set unit-circle} \rangle \langle x \neq i \rangle \text{circline-set-def inversion-unit-circle}$
by fastforce

ultimately

have $?l1 = ?l2$
using unique-circline-set[of x inversion x i]
by blast

hence $0_h \in \text{circline-set } ?l2$
by (metis $\langle a \neq 0_h \rangle \text{poincare-line-circline-set}(1)$)

thus False

using $\langle \neg \text{poincare-collinear } \{0_h, x1, x2\} \rangle$
unfolding poincare-collinear-def
using $\langle \text{poincare-collinear } \{x, x1, x2\} \rangle \langle x1 \neq x2 \rangle \langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle \text{poincare-collinear-def unique-poincare-line}$
by auto

qed

lemma limiting-parallels:

assumes $a \in \text{unit-disc}$ and $x1 \in \text{unit-disc}$ and $x2 \in \text{unit-disc}$ and $\neg \text{poincare-on-line } a \ x1 \ x2$

shows $\exists a1 \in \text{unit-disc}. \exists a2 \in \text{unit-disc}.$

$\neg \text{poincare-on-line } a \ a1 \ a2 \wedge$

$\neg \text{poincare-ray-meets-line } a \ a1 \ x1 \ x2 \wedge \neg \text{poincare-ray-meets-line } a \ a2 \ x1 \ x2 \wedge$

$(\forall a' \in \text{unit-disc}. \text{poincare-in-angle } a' \ a1 \ a \ a2 \longrightarrow \text{poincare-ray-meets-line } a \ a' \ x1 \ x2)$ (is ?P $a \ x1 \ x2$)

proof-

have $\neg \text{poincare-collinear } \{a, x1, x2\}$

using $\langle \neg \text{poincare-on-line } a \ x1 \ x2 \rangle$

unfolding poincare-on-line-def

by simp

have $\forall x1 \ x2. x1 \in \text{unit-disc} \wedge x2 \in \text{unit-disc} \wedge \neg \text{poincare-collinear } \{a, x1, x2\} \longrightarrow ?P \ a \ x1 \ x2$ (is ?Q a)

proof (rule wlog-zero[OF $\langle a \in \text{unit-disc} \rangle$])

fix $a \ u$

assume *: $u \in \text{unit-disc} \text{ cmod } a < 1$

hence $uf: \text{unit-disc-fix } (\text{blaschke } a)$

by simp

assume **: ?Q (moebius-pt (blaschke a) u)

show ?Q u

proof safe

fix $x1 \ x2$

let $?M = \text{moebius-pt } (\text{blaschke } a)$

assume $xx: x1 \in \text{unit-disc} \ x2 \in \text{unit-disc} \neg \text{poincare-collinear } \{u, x1, x2\}$

hence $MM: ?M\ x1 \in \text{unit-disc} \wedge ?M\ x2 \in \text{unit-disc} \wedge \neg \text{poincare-collinear} \{?M\ u, ?M\ x1, ?M\ x2\}$
using *
by *auto*
show $?P\ u\ x1\ x2$ (**is** $\exists a1 \in \text{unit-disc}. \exists a2 \in \text{unit-disc}. ?P'\ a1\ a2\ u\ x1\ x2$)
proof–
obtain $Ma1\ Ma2$ **where** $MM: Ma1 \in \text{unit-disc}\ Ma2 \in \text{unit-disc}\ ?P'\ Ma1\ Ma2\ (?M\ u)\ (?M\ x1)\ (?M\ x2)$
using $**[\text{rule-format}, \text{OF}\ MM]$
by *blast*
hence $MM': \forall a' \in \text{unit-disc}. \text{poincare-in-angle}\ a'\ Ma1\ (?M\ u)\ Ma2 \longrightarrow \text{poincare-ray-meets-line}\ (?M\ u)\ a'\ (?M\ x1)\ (?M\ x2)$
by *auto*
obtain $a1\ a2$ **where** $a: a1 \in \text{unit-disc}\ a2 \in \text{unit-disc}\ ?M\ a1 = Ma1\ ?M\ a2 = Ma2$
using *uf*
by (*metis* $\langle Ma1 \in \text{unit-disc} \rangle \langle Ma2 \in \text{unit-disc} \rangle \text{image-iff}\ \text{unit-disc-fix-iff}$)

have $\forall a' \in \text{unit-disc}. \text{poincare-in-angle}\ a'\ a1\ u\ a2 \longrightarrow \text{poincare-ray-meets-line}\ u\ a'\ x1\ x2$
proof *safe*
fix a'
assume $a' \in \text{unit-disc}\ \text{poincare-in-angle}\ a'\ a1\ u\ a2$
thus $\text{poincare-ray-meets-line}\ u\ a'\ x1\ x2$
using $MM(1-2)\ MM'[\text{rule-format}, \text{of}\ ?M\ a'] * \text{uf}\ a\ xx$
by (*meson* $\text{unit-disc-fix-discI}\ \text{unit-disc-fix-preserves-poincare-in-angle}\ \text{unit-disc-fix-preserves-poincare-ray-meets-line}$)
qed

hence $?P'\ a1\ a2\ u\ x1\ x2$
using $MM * \text{uf}\ xx\ a$
by *auto*

thus *thesis*
using $\langle a1 \in \text{unit-disc} \rangle \langle a2 \in \text{unit-disc} \rangle$
by *blast*
qed
qed
next
show $?Q\ 0_h$
proof *safe*
fix $x1\ x2$
assume $x1 \in \text{unit-disc}\ x2 \in \text{unit-disc}$
assume $\neg \text{poincare-collinear} \{0_h, x1, x2\}$
show $?P\ 0_h\ x1\ x2$
proof–
let $?lx = \text{poincare-line}\ x1\ x2$

have $x1 \neq x2$
using $\langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle \langle \neg \text{poincare-collinear} \{0_h, x1, x2\} \rangle$
using *poincare-collinear3-between*
by *auto*

have $lx: \text{is-poincare-line}\ ?lx$
using *is-poincare-line-poincare-line*[*OF* $\langle x1 \neq x2 \rangle$]
by *simp*

obtain $i1\ i2$ **where** $\text{ideal-points}\ ?lx = \{i1, i2\}$
by (*meson* $\langle x1 \neq x2 \rangle \text{is-poincare-line-poincare-line}\ \text{obtain-ideal-points}$)

let $?li = \text{poincare-line}\ i1\ i2$
let $?i1 = \text{to-complex}\ i1$
let $?i2 = \text{to-complex}\ i2$

have $i1 \in \text{unit-circle-set}\ i2 \in \text{unit-circle-set}$
using $lx\ \langle \text{ideal-points}\ ?lx = \{i1, i2\} \rangle$
unfolding *unit-circle-set-def*
by (*metis* $\text{ideal-points-on-unit-circle}\ \text{insertI1}, \text{metis}\ \text{ideal-points-on-unit-circle}\ \text{insertI1}\ \text{insertI2}$)

have $i1 \neq i2$
using $\langle \text{ideal-points}\ ?lx = \{i1, i2\} \rangle \langle x1 \in \text{unit-disc} \rangle \langle x1 \neq x2 \rangle \langle x2 \in \text{unit-disc} \rangle \text{ideal-points-different}(1)$

by *blast*

let ?a1 = of-complex (?i1 / 2)
let ?a2 = of-complex (?i2 / 2)
let ?la = poincare-line ?a1 ?a2

have ?a1 ∈ unit-disc ?a2 ∈ unit-disc
 using ⟨i1 ∈ unit-circle-set⟩ ⟨i2 ∈ unit-circle-set⟩
 unfolding unit-circle-set-def unit-disc-def disc-def circline-set-def
 by auto (transfer, transfer, case-tac i1, case-tac i2, simp add: vec-cnj-def)+

have ?a1 ≠ 0_h ?a2 ≠ 0_h
 using ⟨i1 ∈ unit-circle-set⟩ ⟨i2 ∈ unit-circle-set⟩
 unfolding unit-circle-set-def
 by auto

have ?a1 ≠ ?a2
 using ⟨i1 ≠ i2⟩
 by (metis ⟨i1 ∈ unit-circle-set⟩ ⟨i2 ∈ unit-circle-set⟩ circline-set-def divide-cancel-right inversion-infty inversion-unit-circle mem-Collect-eq of-complex-to-complex of-complex-zero to-complex-of-complex unit-circle-set-def zero-neq-numeral)

have poincare-collinear {0_h, ?a1, i1}
 unfolding poincare-collinear-def
 using ⟨?a1 ≠ 0_h⟩[symmetric] is-poincare-line-poincare-line[of 0_h ?a1]
 unfolding circline-set-def
 apply (rule-tac x=poincare-line 0_h ?a1 in exI, auto)
 apply (transfer, transfer, auto simp add: vec-cnj-def)
 done

have poincare-collinear {0_h, ?a2, i2}
 unfolding poincare-collinear-def
 using ⟨?a2 ≠ 0_h⟩[symmetric] is-poincare-line-poincare-line[of 0_h ?a2]
 unfolding circline-set-def
 apply (rule-tac x=poincare-line 0_h ?a2 in exI, auto)
 apply (transfer, transfer, auto simp add: vec-cnj-def)
 done

have ¬ poincare-ray-meets-line 0_h ?a1 x1 x2
 using tangent-not-meet[of x1 x2 i1 ?a1]
 using ⟨x1 ∈ unit-disc⟩ ⟨x2 ∈ unit-disc⟩ ⟨?a1 ∈ unit-disc⟩ ⟨x1 ≠ x2⟩ ⟨¬ poincare-collinear {0_h, x1, x2}⟩
 using ⟨ideal-points ?lx = {i1, i2}⟩ ⟨?a1 ≠ 0_h⟩ ⟨poincare-collinear {0_h, ?a1, i1}⟩
 by simp

moreover

have ¬ poincare-ray-meets-line 0_h ?a2 x1 x2
 using tangent-not-meet[of x1 x2 i2 ?a2]
 using ⟨x1 ∈ unit-disc⟩ ⟨x2 ∈ unit-disc⟩ ⟨?a2 ∈ unit-disc⟩ ⟨x1 ≠ x2⟩ ⟨¬ poincare-collinear {0_h, x1, x2}⟩
 using ⟨ideal-points ?lx = {i1, i2}⟩ ⟨?a2 ≠ 0_h⟩ ⟨poincare-collinear {0_h, ?a2, i2}⟩
 by simp

moreover

have ∀ a' ∈ unit-disc. poincare-in-angle a' ?a1 0_h ?a2 → poincare-ray-meets-line 0_h a' x1 x2
 unfolding poincare-in-angle-def

proof safe

fix a' a

assume *: a' ∈ unit-disc a ∈ unit-disc poincare-on-ray a' 0_h a a' ≠ 0_h
 poincare-between ?a1 a ?a2 a ≠ ?a1 a ≠ ?a2

show poincare-ray-meets-line 0_h a' x1 x2

proof −

have ∀ a' a1 a2 x1 x2 i1 i2.

 a' ∈ unit-disc ∧ x1 ∈ unit-disc ∧ x2 ∈ unit-disc ∧ x1 ≠ x2 ∧

 ¬ poincare-collinear {0_h, x1, x2} ∧ ideal-points (poincare-line x1 x2) = {i1, i2} ∧

 a1 = of-complex (to-complex i1 / 2) ∧ a2 = of-complex (to-complex i2 / 2) ∧

 i1 ≠ i2 ∧ a1 ≠ a2 ∧ poincare-collinear {0_h, a1, i1} ∧ poincare-collinear {0_h, a2, i2} ∧

```

    a1 ∈ unit-disc ∧ a2 ∈ unit-disc ∧ i1 ∈ unit-circle-set ∧ i2 ∈ unit-circle-set ∧
    poincare-on-ray a' 0h a ∧ a' ≠ 0h ∧ poincare-between a1 a a2 ∧ a ≠ a1 ∧ a ≠ a2 →
    poincare-ray-meets-line 0h a' x1 x2 (is ∨ a' a1 a2 x1 x2 i1 i2. ?R 0h a' a1 a2 x1 x2 i1 i2 a)
proof (rule wlog-rotation-to-positive-x-axis[OF ⟨a ∈ unit-disc⟩])
let ?R' = λ a zero. ∀ a' a1 a2 x1 x2 i1 i2. ?R zero a' a1 a2 x1 x2 i1 i2 a
fix xa
assume xa: is-real xa 0 < Re xa Re xa < 1
let ?a = of-complex xa
show ?R' ?a 0h
proof safe
  fix a' a1 a2 x1 x2 i1 i2
  let ?i1 = to-complex i1 and ?i2 = to-complex i2
  let ?a1 = of-complex (?i1 / 2) and ?a2 = of-complex (?i2 / 2)
  let ?la = poincare-line ?a1 ?a2 and ?lx = poincare-line x1 x2 and ?li = poincare-line i1 i2
  assume a' ∈ unit-disc x1 ∈ unit-disc x2 ∈ unit-disc x1 ≠ x2
  assume ¬ poincare-collinear {0h, x1, x2} ideal-points ?lx = {i1, i2}
  assume poincare-on-ray a' 0h ?a a' ≠ 0h
  assume poincare-between ?a1 ?a ?a2 ?a ≠ ?a1 ?a ≠ ?a2
  assume i1 ≠ i2 ?a1 ≠ ?a2 poincare-collinear {0h, ?a1, i1} poincare-collinear {0h, ?a2, i2}
  assume ?a1 ∈ unit-disc ?a2 ∈ unit-disc
  assume i1 ∈ unit-circle-set i2 ∈ unit-circle-set
  show poincare-ray-meets-line 0h a' x1 x2
  proof -
    have ?lx = ?li
      using ⟨ideal-points ?lx = {i1, i2}⟩ ⟨x1 ≠ x2⟩ ideal-points-line-unique
      by auto

    have lx: is-poincare-line ?lx
      using is-poincare-line-poincare-line[OF ⟨x1 ≠ x2⟩]
      by simp

    have x1 ∈ circline-set ?lx x2 ∈ circline-set ?lx
      using lx ⟨x1 ≠ x2⟩
      by auto

    have ?lx ≠ x-axis
      using ⟨¬ poincare-collinear {0h, x1, x2}⟩ ⟨x1 ∈ circline-set ?lx⟩ ⟨x2 ∈ circline-set ?lx⟩ lx
      unfolding poincare-collinear-def
      by auto

    have 0h ∉ circline-set ?lx
      using ⟨¬ poincare-collinear {0h, x1, x2}⟩ lx ⟨x1 ∈ circline-set ?lx⟩ ⟨x2 ∈ circline-set ?lx⟩
      unfolding poincare-collinear-def
      by auto

    have xa ≠ 0 ?a ≠ 0h
      using xa
      by auto
    hence 0h ≠ ?a
      by metis

    have ?a ∈ positive-x-axis
      using xa
      unfolding positive-x-axis-def
      by simp

    have ?a ∈ unit-disc
      using xa
      by (auto simp add: cmod-eq-Re)

    have ?a ∈ circline-set ?la
      using ⟨poincare-between ?a1 ?a ?a2⟩
    using ⟨?a1 ≠ ?a2⟩ ⟨?a ∈ unit-disc⟩ ⟨?a1 ∈ unit-disc⟩ ⟨?a2 ∈ unit-disc⟩ poincare-between-poincare-line-uzv
      by blast

```

```

have ?a1 ∈ circline-set ?la ?a2 ∈ circline-set ?la
  by (auto simp add: ⟨?a1 ≠ ?a2⟩)

have la: is-poincare-line ?la
  using is-poincare-line-poincare-line[OF ⟨?a1 ≠ ?a2⟩]
  by simp

have inv: inversion i1 = i1 inversion i2 = i2
  using ⟨i1 ∈ unit-circle-set⟩ ⟨i2 ∈ unit-circle-set⟩
  by (auto simp add: circline-set-def unit-circle-set-def)

have i1 ≠ ∞h i2 ≠ ∞h
  using inv
  by auto

have ?a1 ∉ circline-set x-axis ∧ ?a2 ∉ circline-set x-axis
proof (rule ccontr)
  assume ¬ ?thesis
  hence ?a1 ∈ circline-set x-axis ∨ ?a2 ∈ circline-set x-axis
    by auto
  hence ?la = x-axis
  proof
    assume ?a1 ∈ circline-set x-axis
    hence {?a, ?a1} ⊆ circline-set ?la ∩ circline-set x-axis
      using ⟨?a ∈ circline-set ?la⟩ ⟨?a1 ∈ circline-set ?la⟩ ⟨?a ∈ positive-x-axis⟩
      using circline-set-x-axis-I xa(1)
      by blast
    thus ?la = x-axis
      using unique-is-poincare-line[of ?a ?a1 ?la x-axis]
      using ⟨?a1 ∈ unit-disc⟩ ⟨?a ∈ unit-disc⟩ la ⟨?a ≠ ?a1⟩
      by auto
  next
    assume ?a2 ∈ circline-set x-axis
    hence {?a, ?a2} ⊆ circline-set ?la ∩ circline-set x-axis
      using ⟨?a ∈ circline-set ?la⟩ ⟨?a2 ∈ circline-set ?la⟩ ⟨?a ∈ positive-x-axis⟩
      using circline-set-x-axis-I xa(1)
      by blast
    thus ?la = x-axis
      using unique-is-poincare-line[of ?a ?a2 ?la x-axis]
      using ⟨?a2 ∈ unit-disc⟩ ⟨?a ∈ unit-disc⟩ la ⟨?a ≠ ?a2⟩
      by auto
  qed

  hence i1 ∈ circline-set x-axis ∧ i2 ∈ circline-set x-axis
    using ⟨?a1 ∈ circline-set ?la⟩ ⟨?a2 ∈ circline-set ?la⟩
    by (metis ⟨i1 ≠ ∞h⟩ ⟨i2 ≠ ∞h⟩ ⟨of-complex (to-complex i1 / 2) ∈ unit-disc⟩ ⟨of-complex (to-complex i2 / 2) ∈ unit-disc⟩ ⟨poincare-collinear {0h, of-complex (to-complex i1 / 2), i1}⟩ ⟨poincare-collinear {0h, of-complex (to-complex i2 / 2), i2}⟩ divide-eq-0-iff inf-not-of-complex inv(1) inv(2) inversion-not-eq-unit-disc of-complex-to-complex of-complex-zero-iff poincare-collinear3-poincare-lines-equal-general poincare-line-0-real-is-x-axis poincare-line-circline-set(2) zero-in-unit-disc zero-neq-numeral)

  thus False
    using ⟨?lx ≠ x-axis⟩ unique-is-poincare-line-general[of i1 i2 ?li x-axis] ⟨i1 ≠ i2⟩ inv ⟨?lx = ?li⟩
    by auto
  qed

  hence ?la ≠ x-axis
    using ⟨?a1 ≠ ?a2⟩ poincare-line-circline-set(1)
    by fastforce

  have intersects-x-axis-positive ?la
    using intersects-x-axis-positive-iff[of ?la] ⟨?la ≠ x-axis⟩ ⟨?a ∈ circline-set ?la⟩ la
    using ⟨?a ∈ unit-disc⟩ ⟨?a ∈ positive-x-axis⟩
    by auto

  have intersects-x-axis ?lx

```

proof-

have $\text{Arg}(\text{to-complex } ?a1) * \text{Arg}(\text{to-complex } ?a2) < 0$
using $\langle \text{poincare-between } ?a1 ?a ?a2 \rangle \langle ?a1 \in \text{unit-disc} \rangle \langle ?a2 \in \text{unit-disc} \rangle$
using $\text{poincare-between-x-axis-intersection}[\text{of } ?a1 ?a2 \text{ of-complex } xa]$
using $\langle ?a1 \neq ?a2 \rangle \langle ?a \in \text{unit-disc} \rangle \langle ?a1 \notin \text{circline-set } x\text{-axis} \wedge ?a2 \notin \text{circline-set } x\text{-axis} \rangle \langle ?a \in \text{positive-x-axis} \rangle$
using $\langle ?a \in \text{circline-set } ?la \rangle$
unfolding $\text{positive-x-axis-def}$
by simp

moreover

have $\bigwedge x y x' y' :: \text{real. } [\text{sgn } x' = \text{sgn } x; \text{sgn } y' = \text{sgn } y] \implies x*y < 0 \longleftrightarrow x'*y' < 0$
by $(\text{metis } \text{sgn-less } \text{sgn-mult})$

ultimately

have $\text{Im}(\text{to-complex } ?a1) * \text{Im}(\text{to-complex } ?a2) < 0$
using $\text{arg-Im-sgn}[\text{of } \text{to-complex } ?a1] \text{ arg-Im-sgn}[\text{of } \text{to-complex } ?a2]$
using $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a2 \in \text{unit-disc} \rangle \langle ?a1 \notin \text{circline-set } x\text{-axis} \wedge ?a2 \notin \text{circline-set } x\text{-axis} \rangle$
using $\text{inf-or-of-complex}[\text{of } ?a1] \text{ inf-or-of-complex}[\text{of } ?a2] \text{ circline-set-x-axis}$
by $(\text{metis } \text{circline-set-x-axis-I } \text{to-complex-of-complex})$

thus $?thesis$

using $\text{ideal-points-intersects-x-axis}[\text{of } ?lx \ i1 \ i2]$
using $\langle \text{ideal-points } ?lx = \{i1, i2\} \rangle \langle ?lx \neq x\text{-axis} \rangle$
by simp

qed

have $\text{intersects-x-axis-positive } ?lx$

proof-

have $\text{cmod } ?i1 = 1 \ \text{cmod } ?i2 = 1$
using $\langle i1 \in \text{unit-circle-set} \rangle \langle i2 \in \text{unit-circle-set} \rangle$
unfolding $\text{unit-circle-set-def}$
by auto

let $?a1' = ?i1 / 2$ **and** $?a2' = ?i2 / 2$

let $?Aa1 = i * (?a1' * \text{cnj } ?a2' - ?a2' * \text{cnj } ?a1')$ **and**
 $?Ba1 = i * (?a2' * \text{cor } ((\text{cmod } ?a1')^2 + 1) - ?a1' * \text{cor } ((\text{cmod } ?a2')^2 + 1))$

have $?Aa1 \neq 0 \vee ?Ba1 \neq 0$

using $\langle \text{cmod } (\text{to-complex } i1) = 1 \rangle \langle \text{cmod } (\text{to-complex } i2) = 1 \rangle \langle ?a1 \neq ?a2 \rangle$
by $(\text{auto } \text{simp } \text{add: } \text{power-divide } \text{complex-mult-cnj-cmod})$

have $\text{is-real } ?Aa1$

by simp

have $?a1 \neq \text{inversion } ?a2$

using $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a2 \in \text{unit-disc} \rangle \text{inversion-noteq-unit-disc}$ **by** fastforce

hence $\text{Re } ?Ba1 / \text{Re } ?Aa1 < -1$

using $\langle \text{intersects-x-axis-positive } ?la \rangle \langle ?a1 \neq ?a2 \rangle$
using $\text{intersects-x-axis-positive-mk-circline}[\text{of } ?Aa1 ?Ba1] \langle ?Aa1 \neq 0 \vee ?Ba1 \neq 0 \rangle \langle \text{is-real } ?Aa1 \rangle$
using $\text{poincare-line-non-homogenous}[\text{of } ?a1 ?a2]$
by $(\text{simp } \text{add: } \text{Let-def})$

moreover

let $?i1' = \text{to-complex } i1$ **and** $?i2' = \text{to-complex } i2$

let $?Ai1 = i * (?i1' * \text{cnj } ?i2' - ?i2' * \text{cnj } ?i1')$ **and**
 $?Bi1 = i * (?i2' * \text{cor } ((\text{cmod } ?i1')^2 + 1) - ?i1' * \text{cor } ((\text{cmod } ?i2')^2 + 1))$

have $?Ai1 \neq 0 \vee ?Bi1 \neq 0$

using $\langle \text{cmod } (\text{to-complex } i1) = 1 \rangle \langle \text{cmod } (\text{to-complex } i2) = 1 \rangle \langle ?a1 \neq ?a2 \rangle$
by $(\text{auto } \text{simp } \text{add: } \text{power-divide } \text{complex-mult-cnj-cmod})$

```

have is-real ?Ai1
  by simp

have sgn (Re ?Bi1 / Re ?Ai1) = sgn (Re ?Ba1 / Re ?Aa1)
proof-
  have Re ?Bi1 / Re ?Ai1 = (Im ?i1 * 2 - Im ?i2 * 2) /
    (Im ?i2 * (Re ?i1 * 2) - Im ?i1 * (Re ?i2 * 2))
    using ⟨cmod ?i1 = 1⟩ ⟨cmod ?i2 = 1⟩
    by (auto simp add: complex-mult-cnj-cmod field-simps)
  also have ... = (Im ?i1 - Im ?i2) /
    (Im ?i2 * (Re ?i1) - Im ?i1 * (Re ?i2)) (is ... = ?expr)
    apply (subst left-diff-distrib[symmetric])
    apply (subst semiring-normalization-rules(18))+
    apply (subst left-diff-distrib[symmetric])
    by (metis mult.commute mult-divide-mult-cancel-left-if zero-neq-numeral)
  finally have 1: Re ?Bi1 / Re ?Ai1 = (Im ?i1 - Im ?i2) / (Im ?i2 * (Re ?i1) - Im ?i1 * (Re ?i2))
  .

```

```

have Re ?Ba1 / Re ?Aa1 = (Im ?i1 * 20 - Im ?i2 * 20) /
  (Im ?i2 * (Re ?i1 * 16) - Im ?i1 * (Re ?i2 * 16))
  using ⟨cmod (to-complex i1) = 1⟩ ⟨cmod (to-complex i2) = 1⟩
  by (auto simp add: complex-mult-cnj-cmod field-simps)
also have ... = (20 / 16) * ((Im ?i1 - Im ?i2) /
  (Im ?i2 * (Re ?i1) - Im ?i1 * (Re ?i2)))
  apply (subst left-diff-distrib[symmetric])+
  apply (subst semiring-normalization-rules(18))+
  apply (subst left-diff-distrib[symmetric])+
by (metis (no-types, opaque-lifting) field-class.field-divide-inverse mult.commute times-divide-times-eq)
finally have 2: Re ?Ba1 / Re ?Aa1 = (5 / 4) * ((Im ?i1 - Im ?i2) / (Im ?i2 * (Re ?i1) - Im ?i1
* (Re ?i2)))
  by simp

have ?expr ≠ 0
  using ⟨Re ?Ba1 / Re ?Aa1 < -1⟩
  apply (subst (asm) 2)
  by linarith
thus ?thesis
  apply (subst 1, subst 2)
  apply (simp only: sgn-mult)
  by simp
qed

```

moreover

```

have i1 ≠ inversion i2
  by (simp add: ⟨i1 ≠ i2⟩ inv(2))

```

```

have (Re ?Bi1 / Re ?Ai1)2 > 1

```

```

proof-
  have ?Ai1 = 0 ∨ (Re ?Bi1)2 > (Re ?Ai1)2
    using ⟨intersects-x-axis ?lx⟩
    using ⟨i1 ≠ i2⟩ ⟨i1 ≠ ∞h⟩ ⟨i2 ≠ ∞h⟩ ⟨i1 ≠ inversion i2⟩
    using intersects-x-axis-mk-circline[of ?Ai1 ?Bi1] ⟨?Ai1 ≠ 0 ∨ ?Bi1 ≠ 0⟩ ⟨is-real ?Ai1⟩
    using poincare-line-non-homogenous[of i1 i2] ⟨?lx = ?li⟩
    by metis

```

moreover

```

have ?Ai1 ≠ 0
proof (rule ccontr)
  assume ¬ ?thesis
  hence 0h ∈ circline-set ?li
    unfolding circline-set-def
  apply simp
  apply (transfer, transfer, case-tac i1, case-tac i2)

```

```

    by (auto simp add: vec-cnj-def field-simps)
  thus False
    using ⟨0h ∉ circline-set ?lx⟩ ⟨?lx = ?li⟩
    by simp
qed

ultimately

have (Re ?Bi1)2 > (Re ?Ai1)2
  by auto

moreover

have Re ?Ai1 ≠ 0
  using ⟨is-real ?Ai1⟩ ⟨?Ai1 ≠ 0⟩
  by (simp add: complex-eq-iff)

ultimately

show ?thesis
  by (simp add: power-divide)
qed

moreover

{
  fix x1 x2 :: real
  assume sgn x1 = sgn x2 x1 < -1 x22 > 1
  hence x2 < -1
    by (smt one-power2 real-sqrt-abs real-sqrt-less-iff sgn-neg sgn-pos)
}

ultimately

have Re ?Bi1 / Re ?Ai1 < -1
  by metis

thus ?thesis
  using ⟨i1 ≠ i2⟩ ⟨i1 ≠ ∞h⟩ ⟨i2 ≠ ∞h⟩ ⟨i1 ≠ inversion i2⟩
  using intersects-x-axis-positive-mk-circline[of ?Ai1 ?Bi1] ⟨?Ai1 ≠ 0 ∨ ?Bi1 ≠ 0⟩ ⟨is-real ?Ai1⟩
  using poincare-line-non-homogenous[of i1 i2] ⟨?lx = ?li⟩
  by (simp add: Let-def)
qed

then obtain x where x: x ∈ unit-disc x ∈ circline-set ?lx ∩ positive-x-axis
  using intersects-x-axis-positive-iff[OF lx ⟨?lx ≠ x-axis⟩]
  by auto

have poincare-on-ray x 0h a' ∧ poincare-collinear {x1, x2, x}
proof
  show poincare-collinear {x1, x2, x}
    using x lx ⟨x1 ∈ circline-set ?lx⟩ ⟨x2 ∈ circline-set ?lx⟩
    unfolding poincare-collinear-def
    by auto
next
  show poincare-on-ray x 0h a'
    unfolding poincare-on-ray-def
  proof-
    have a' ∈ circline-set x-axis
      using ⟨poincare-on-ray a' 0h ?a⟩ xa ⟨0h ≠ ?a⟩ ⟨xa ≠ 0⟩ ⟨a' ∈ unit-disc⟩
      unfolding poincare-on-ray-def
      using poincare-line-0-real-is-x-axis[of of-complex xa]
      using poincare-between-poincare-line-uvw[of 0h of-complex xa a']
      using poincare-between-poincare-line-uzv[of 0h of-complex xa a']
      by (auto simp add: cmod-eq-Re)

```

then obtain xa' **where** xa' : $a' = \text{of-complex } xa' \text{ is-real } xa'$
using $\langle a' \in \text{unit-disc} \rangle$
using $\text{circline-set-def on-circline-x-axis}$
by auto

hence $-1 < \text{Re } xa' \text{ Re } xa' < 1 \text{ } xa' \neq 0$
using $\langle a' \in \text{unit-disc} \rangle \langle a' \neq 0_h \rangle$
by $(\text{auto simp add: cmod-eq-Re})$

hence $\text{Re } xa' > 0 \text{ Re } xa' < 1 \text{ is-real } xa'$
using $\langle \text{poincare-on-ray } a' 0_h \text{ (of-complex } xa) \rangle$
using $\text{poincare-between-x-axis-0uv}[of \text{ Re } xa' \text{ Re } xa]$
using $\text{poincare-between-x-axis-0uv}[of \text{ Re } xa \text{ Re } xa']$
using $\text{circline-set-positive-x-axis-I}[of \text{ Re } xa']$
using $xa \text{ } xa' \text{ complex-of-real-Re}$
unfolding $\text{poincare-on-ray-def}$
by $(\text{smt of-real-0, linarith, blast})$

moreover

obtain xx **where** $\text{is-real } xx \text{ Re } xx > 0 \text{ Re } xx < 1 \text{ } x = \text{of-complex } xx$
using x
unfolding $\text{positive-x-axis-def}$
using $\text{circline-set-def cmod-eq-Re on-circline-x-axis}$
by auto

ultimately

show $\text{poincare-between } 0_h \text{ } x \text{ } a' \vee \text{poincare-between } 0_h \text{ } a' \text{ } x$
using $\langle a' = \text{of-complex } xa' \rangle$
by $(\text{smt } \langle a' \in \text{unit-disc} \rangle \text{ arg-0-iff poincare-between-0uv poincare-between-def to-complex-of-complex})$

$x(1))$

qed

qed

thus $?thesis$
using $\langle x \in \text{unit-disc} \rangle$
unfolding $\text{poincare-ray-meets-line-def poincare-on-line-def}$
by $(\text{metis insert-commute})$

qed

qed

next

show $a \neq 0_h$

proof (rule ccontr)

assume $\neg ?thesis$

then obtain k **where** $k < 0 \text{ to-complex } ?a1 = \text{cor } k * \text{to-complex } ?a2$

using $\text{poincare-between-u0v}[OF \langle ?a1 \in \text{unit-disc} \rangle \langle ?a2 \in \text{unit-disc} \rangle \langle ?a1 \neq 0_h \rangle \langle ?a2 \neq 0_h \rangle]$

using $\langle \text{poincare-between } ?a1 \text{ } a \text{ } ?a2 \rangle$

by auto

hence $\text{to-complex } i1 = \text{cor } k * \text{to-complex } i2 \text{ } k < 0$

by auto

hence $0_h \in \text{circline-set } (\text{poincare-line } x1 \text{ } x2)$

using $\text{ideal-points-proportional}[of \text{poincare-line } x1 \text{ } x2 \text{ } i1 \text{ } i2 \text{ } k] \langle \text{ideal-points } (\text{poincare-line } x1 \text{ } x2) = \{i1,$

$i2\} \rangle$

using $\text{is-poincare-line-poincare-line}[OF \langle x1 \neq x2 \rangle]$

by simp

thus False

using $\langle \neg \text{poincare-collinear } \{0_h, x1, x2\} \rangle$

using $\text{is-poincare-line-poincare-line}[OF \langle x1 \neq x2 \rangle]$

unfolding $\text{poincare-collinear-def}$

by $(\text{meson } \langle x1 \neq x2 \rangle \text{ empty-subsetI insert-subset poincare-line-circline-set}(1) \text{ poincare-line-circline-set}(2))$

qed

next

fix $\varphi \text{ } u$

let $?R' = \lambda a \text{ zero. } \forall a' a1 a2 \text{ } x1 \text{ } x2 \text{ } i1 \text{ } i2. ?R \text{ zero } a' a1 a2 \text{ } x1 \text{ } x2 \text{ } i1 \text{ } i2 \text{ } a$


```

let ?M = moebius-pt (moebius-rotation  $\varphi$ )
assume *:  $u \in \text{unit-disc } u \neq 0_h$  and **:  $?R' (?M u) 0_h$ 
have uf:  $\text{unit-disc-fix (moebius-rotation } \varphi)$ 
  by simp
have ?M  $0_h = 0_h$ 
  by auto
hence **:  $?R' (?M u) (?M 0_h)$ 
  using **
  by simp
show  $?R' u 0_h$ 
proof (rule allI)+
  fix a' a1 a2 x1 x2 i1 i2
  have i1:  $i1 \in \text{unit-circle-set} \longrightarrow \text{moebius-pt (moebius-rotation } \varphi) (\text{of-complex (to-complex } i1 / 2)) =$ 
of-complex (to-complex (moebius-pt (moebius-rotation  $\varphi$ ) i1) / 2)
  using unit-circle-set-def by force

  have i2:  $i2 \in \text{unit-circle-set} \longrightarrow \text{moebius-pt (moebius-rotation } \varphi) (\text{of-complex (to-complex } i2 / 2)) =$ 
of-complex (to-complex (moebius-pt (moebius-rotation  $\varphi$ ) i2) / 2)
  using unit-circle-set-def by force

  show  $?R 0_h a' a1 a2 x1 x2 i1 i2 u$ 
  using **[rule-format, of ?M a' ?M x1 ?M x2 ?M i1 ?M i2 ?M a1 ?M a2] uf *
  apply (auto simp del: moebius-pt-moebius-rotation-zero moebius-pt-moebius-rotation)
  using i1 i2
  by simp
qed
qed
thus ?thesis
  using  $\langle a' \in \text{unit-disc} \rangle \langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle \langle x1 \neq x2 \rangle$ 
  using  $\langle \neg \text{poincare-collinear } \{0_h, x1, x2\} \rangle \langle \text{ideal-points } ?lx = \{i1, i2\} \rangle \langle i1 \neq i2 \rangle$ 
  using  $\langle ?a1 \neq ?a2 \rangle \langle \text{poincare-collinear } \{0_h, ?a1, i1\} \rangle \langle \text{poincare-collinear } \{0_h, ?a2, i2\} \rangle$ 
  using  $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a2 \in \text{unit-disc} \rangle \langle i1 \in \text{unit-circle-set} \rangle \langle i2 \in \text{unit-circle-set} \rangle$ 
  using  $\langle \text{poincare-on-ray } a' 0_h a \rangle \langle a' \neq 0_h \rangle \langle \text{poincare-between } ?a1 a ?a2 \rangle \langle a \neq ?a1 \rangle \langle a \neq ?a2 \rangle$ 
  by blast
qed
qed

moreover

have  $\neg \text{poincare-on-line } 0_h ?a1 ?a2$ 
proof
  assume *:  $\text{poincare-on-line } 0_h ?a1 ?a2$ 
  hence  $\text{poincare-collinear } \{0_h, ?a1, ?a2\}$ 
  unfolding  $\text{poincare-on-line-def}$ 
  by simp
  hence  $\text{poincare-line } 0_h ?a1 = \text{poincare-line } 0_h ?a2$ 
  using  $\text{poincare-collinear3-poincare-lines-equal-general}[of 0_h ?a1 ?a2]$ 
  using  $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a1 \neq 0_h \rangle \langle ?a2 \in \text{unit-disc} \rangle \langle ?a2 \neq 0_h \rangle$ 
  by (metis inversion-noteq-unit-disc zero-in-unit-disc)

  have  $i1 \in \text{circline-set (poincare-line } 0_h ?a1)$ 
  using  $\langle \text{poincare-collinear } \{0_h, ?a1, i1\} \rangle$ 
  using  $\text{poincare-collinear3-poincare-line-general}[of i1 0_h ?a1]$ 
  using  $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a1 \neq 0_h \rangle$ 
  by (metis insert-commute inversion-noteq-unit-disc zero-in-unit-disc)
  moreover
  have  $i2 \in \text{circline-set (poincare-line } 0_h ?a1)$ 
  using  $\langle \text{poincare-collinear } \{0_h, ?a2, i2\} \rangle$ 
  using  $\text{poincare-collinear3-poincare-line-general}[of i2 0_h ?a2]$ 
  using  $\langle ?a2 \in \text{unit-disc} \rangle \langle ?a2 \neq 0_h \rangle \langle \text{poincare-line } 0_h ?a1 = \text{poincare-line } 0_h ?a2 \rangle$ 
  by (metis insert-commute inversion-noteq-unit-disc zero-in-unit-disc)

ultimately

have  $\text{poincare-collinear } \{0_h, i1, i2\}$ 
  using  $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a1 \neq 0_h \rangle \langle \text{poincare-collinear } \{0_h, ?a1, i1\} \rangle$ 

```

```

    by (smt insert-subset poincare-collinear-def unique-poincare-line zero-in-unit-disc)
  hence  $0_h \in \text{circline-set } (\text{poincare-line } i1 \ i2)$ 
    using poincare-collinear3-poincare-line-general[of  $0_h \ i1 \ i2$ ]
    using  $\langle i1 \neq i2 \rangle \langle i2 \in \text{unit-circle-set} \rangle \text{unit-circle-set-def}$ 
    by force

  moreover

  have  $?lx = ?li$ 
    using  $\langle \text{ideal-points } ?lx = \{i1, i2\} \rangle \langle x1 \neq x2 \rangle \text{ideal-points-line-unique}$ 
    by auto

  ultimately

  show False
    using  $\langle \neg \text{poincare-collinear } \{0_h, x1, x2\} \rangle$ 
    using  $\langle x1 \neq x2 \rangle \text{poincare-line-poincare-collinear3-general}$ 
    by auto
  qed

  ultimately

  show ?thesis
    using  $\langle ?a1 \in \text{unit-disc} \rangle \langle ?a2 \in \text{unit-disc} \rangle$ 
    by blast
  qed
qed
qed
qed
thus ?thesis
  using  $\langle x1 \in \text{unit-disc} \rangle \langle x2 \in \text{unit-disc} \rangle \langle \neg \text{poincare-collinear } \{a, x1, x2\} \rangle$ 
  by blast
qed

```

11.9 Interpretation of locales

global-interpretation *PoincareTarskiAbsolute*: *TarskiAbsolute* **where** *cong* = *p-congruent* **and** *betw* = *p-between*

defines *p-on-line* = *PoincareTarskiAbsolute.on-line* **and**
p-on-ray = *PoincareTarskiAbsolute.on-ray* **and**
p-in-angle = *PoincareTarskiAbsolute.in-angle* **and**
p-ray-meets-line = *PoincareTarskiAbsolute.ray-meets-line*

proof–

show *TarskiAbsolute* *p-congruent* *p-between*

proof

1. Reflexivity of congruence

```

  fix  $x \ y$ 
  show p-congruent  $x \ y \ y \ x$ 
    unfolding p-congruent-def
    by transfer (simp add: poincare-distance-sym)
  next

```

2. Transitivity of congruence

```

  fix  $x \ y \ z \ u \ v \ w$ 
  show p-congruent  $x \ y \ z \ u \ \wedge \ \text{p-congruent } x \ y \ v \ w \longrightarrow \text{p-congruent } z \ u \ v \ w$ 
    by (transfer, simp)
  next

```

3. Identity of congruence

```

  fix  $x \ y \ z$ 
  show p-congruent  $x \ y \ z \ z \longrightarrow x = y$ 
    unfolding p-congruent-def
    by transfer (simp add: poincare-distance-eq-0-iff)
  next

```

4. Segment construction

```

fix x y a b
show  $\exists z. p\text{-between } x y z \wedge p\text{-congruent } y z a b$ 
  using segment-construction
  unfolding p-congruent-def
  by transfer (simp, blast)
next

```

5. Five segment

```

fix x y z x' y' z' u u'
show  $x \neq y \wedge p\text{-between } x y z \wedge p\text{-between } x' y' z' \wedge$ 
   $p\text{-congruent } x y x' y' \wedge p\text{-congruent } y z y' z' \wedge$ 
   $p\text{-congruent } x u x' u' \wedge p\text{-congruent } y u y' u' \longrightarrow$ 
   $p\text{-congruent } z u z' u'$ 
  unfolding p-congruent-def
  apply transfer
  using five-segment-axiom
  by meson
next

```

6. Identity of betweenness

```

fix x y
show  $p\text{-between } x y x \longrightarrow x = y$ 
  by transfer (simp add: poincare-between-sum-distances poincare-distance-eq-0-iff poincare-distance-sym)
next

```

7. Pasch

```

fix x y z u v
show  $p\text{-between } x u z \wedge p\text{-between } y v z \longrightarrow (\exists a. p\text{-between } u a y \wedge p\text{-between } x a v)$ 
  apply transfer
  using Pasch
  by blast
next

```

8. Lower dimension

```

show  $\exists a. \exists b. \exists c. \neg p\text{-between } a b c \wedge \neg p\text{-between } b c a \wedge \neg p\text{-between } c a b$ 
  apply (transfer)
  using lower-dimension-axiom
  by simp
next

```

9. Upper dimension

```

fix x y z u v
show  $p\text{-congruent } x u x v \wedge p\text{-congruent } y u y v \wedge p\text{-congruent } z u z v \wedge u \neq v \longrightarrow$ 
   $p\text{-between } x y z \vee p\text{-between } y z x \vee p\text{-between } z x y$ 
  unfolding p-congruent-def
  by (transfer, simp add: upper-dimension-axiom)
qed
qed

```

interpretation *PoincareTarskiHyperbolic: TarskiHyperbolic*
where *cong* = *p-congruent* **and** *betw* = *p-between*
proof

10. Euclid negation

```

show  $\exists a b c d t. p\text{-between } a d t \wedge p\text{-between } b d c \wedge a \neq d \wedge$ 
   $(\forall x y. p\text{-between } a b x \wedge p\text{-between } a c y \longrightarrow \neg p\text{-between } x t y)$ 
  using negated-euclidean-axiom
  by transfer (auto, blast)
next
fix a x1 x2
assume  $\neg \text{TarskiAbsolute.on-line } p\text{-between } a x1 x2$ 
hence  $\neg p\text{-on-line } a x1 x2$ 
  using TarskiAbsolute.on-line-def[OF PoincareTarskiAbsolute.TarskiAbsolute-axioms]

```

using *PoincareTarskiAbsolute.on-line-def*
 by *simp*

11. Limiting parallels

thus $\exists a1\ a2.$
 \neg *TarskiAbsolute.on-line* *p-between* *a* *a1* *a2* \wedge
 \neg *TarskiAbsolute.ray-meets-line* *p-between* *a* *a1* *x1* *x2* \wedge
 \neg *TarskiAbsolute.ray-meets-line* *p-between* *a* *a2* *x1* *x2* \wedge
 $(\forall a'. \textit{TarskiAbsolute.in-angle } p\text{-between } a' a1 a2 \longrightarrow \textit{TarskiAbsolute.ray-meets-line } p\text{-between } a a' x1 x2)$

unfolding *TarskiAbsolute.in-angle-def*[*OF PoincareTarskiAbsolute.TarskiAbsolute-axioms*]
 unfolding *TarskiAbsolute.on-ray-def*[*OF PoincareTarskiAbsolute.TarskiAbsolute-axioms*]
 unfolding *TarskiAbsolute.ray-meets-line-def*[*OF PoincareTarskiAbsolute.TarskiAbsolute-axioms*]
 unfolding *TarskiAbsolute.on-ray-def*[*OF PoincareTarskiAbsolute.TarskiAbsolute-axioms*]
 unfolding *TarskiAbsolute.on-line-def*[*OF PoincareTarskiAbsolute.TarskiAbsolute-axioms*]
 unfolding *PoincareTarskiAbsolute.on-line-def*
 apply *transfer*

proof—
 fix *a x1 x2*
 assume *: *a* \in *unit-disc* *x1* \in *unit-disc* *x2* \in *unit-disc*
 \neg (*poincare-between* *a* *x1* *x2* \vee *poincare-between* *x1* *a* *x2* \vee *poincare-between* *x1* *x2* *a*)

hence \neg *poincare-on-line* *a* *x1* *x2*
 using *poincare-collinear3-iff*[*of a x1 x2*]
 using *poincare-between-rev* *poincare-on-line-def* by *blast*

hence $\exists a1 \in \textit{unit-disc}.$
 $\exists a2 \in \textit{unit-disc}.$
 \neg *poincare-on-line* *a* *a1* *a2* \wedge
 \neg *poincare-ray-meets-line* *a* *a1* *x1* *x2* \wedge
 \neg *poincare-ray-meets-line* *a* *a2* *x1* *x2* \wedge
 $(\forall a' \in \textit{unit-disc}.$
 $\textit{poincare-in-angle } a' a1 a2 \longrightarrow$
 $\textit{poincare-ray-meets-line } a a' x1 x2)$

using *limiting-parallels*[*of a x1 x2*] *
 by *blast*

then obtain *a1 a2* where *: *a1* \in *unit-disc* *a2* \in *unit-disc* \neg *poincare-on-line* *a* *a1* *a2*
 \neg *poincare-ray-meets-line* *a* *a2* *x1* *x2*
 \neg *poincare-ray-meets-line* *a* *a1* *x1* *x2*
 $\forall a' \in \textit{unit-disc}.$
 $\textit{poincare-in-angle } a' a1 a2 \longrightarrow$
 $\textit{poincare-ray-meets-line } a a' x1 x2$

by *blast*

have \neg ($\exists x \in \{z. z \in \textit{unit-disc}\}.$
 $(\textit{poincare-between } a\ x\ a1\ \vee$
 $\textit{poincare-between } a\ a1\ x) \wedge$
 $(\textit{poincare-between } x\ x1\ x2\ \vee$
 $\textit{poincare-between } x1\ x\ x2\ \vee$
 $\textit{poincare-between } x1\ x2\ x))$

using $\langle \neg \textit{poincare-ray-meets-line } a\ a1\ x1\ x2 \rangle$
 unfolding *poincare-on-line-def* *poincare-ray-meets-line-def* *poincare-on-ray-def*
 using *poincare-collinear3-iff*[*of - x1 x2*] *poincare-between-rev* *(2, 3)
 by *auto*

moreover

have \neg ($\exists x \in \{z. z \in \textit{unit-disc}\}.$
 $(\textit{poincare-between } a\ x\ a2\ \vee$
 $\textit{poincare-between } a\ a2\ x) \wedge$
 $(\textit{poincare-between } x\ x1\ x2\ \vee$
 $\textit{poincare-between } x1\ x\ x2\ \vee$
 $\textit{poincare-between } x1\ x2\ x))$

using $\langle \neg \textit{poincare-ray-meets-line } a\ a2\ x1\ x2 \rangle$
 unfolding *poincare-on-line-def* *poincare-ray-meets-line-def* *poincare-on-ray-def*
 using *poincare-collinear3-iff*[*of - x1 x2*] *poincare-between-rev* *(2, 3)
 by *auto*

moreover

have \neg (*poincare-between* *a* *a1* *a2* \vee *poincare-between* *a1* *a* *a2* \vee *poincare-between* *a1* *a2* *a*)
 using $\langle \neg \textit{poincare-on-line } a\ a1\ a2 \rangle$ *poincare-collinear3-iff*[*of a a1 a2*]
 using *(1) *(1-2)
 unfolding *poincare-on-line-def*

by simp
moreover
have $(\forall a' \in \{z. z \in \text{unit-disc}\}.$
 $a \neq a1 \wedge$
 $a \neq a2 \wedge$
 $a' \neq a \wedge$
 $(\exists x \in \{z. z \in \text{unit-disc}\}.$
 $\text{poincare-between } a1 \ x \ a2 \wedge$
 $x \neq a1 \wedge$
 $x \neq a2 \wedge$
 $(\text{poincare-between } a \ a' \ x \vee$
 $\text{poincare-between } a \ x \ a')) \longrightarrow$
 $(\exists x \in \{z. z \in \text{unit-disc}\}.$
 $(\text{poincare-between } a \ x \ a' \vee$
 $\text{poincare-between } a \ a' \ x) \wedge$
 $(\text{poincare-between } x \ x1 \ x2 \vee$
 $\text{poincare-between } x1 \ x \ x2 \vee$
 $\text{poincare-between } x1 \ x2 \ x)))$
using $**(6)$
unfolding *poincare-on-line-def* *poincare-in-angle-def* *poincare-ray-meets-line-def* *poincare-on-ray-def*
using *poincare-collinear3-iff*[*of - x1 x2*] *poincare-between-rev* $*(2, 3)$
by auto
ultimately
show $\exists a1 \in \{z. z \in \text{unit-disc}\}.$
 $\exists a2 \in \{z. z \in \text{unit-disc}\}.$
 $\neg (\text{poincare-between } a \ a1 \ a2 \vee \text{poincare-between } a1 \ a \ a2 \vee \text{poincare-between } a1 \ a2 \ a) \wedge$
 $\neg (\exists x \in \{z. z \in \text{unit-disc}\}.$
 $(\text{poincare-between } a \ x \ a1 \vee$
 $\text{poincare-between } a \ a1 \ x) \wedge$
 $(\text{poincare-between } x \ x1 \ x2 \vee$
 $\text{poincare-between } x1 \ x \ x2 \vee$
 $\text{poincare-between } x1 \ x2 \ x)) \wedge$
 $\neg (\exists x \in \{z. z \in \text{unit-disc}\}.$
 $(\text{poincare-between } a \ x \ a2 \vee$
 $\text{poincare-between } a \ a2 \ x) \wedge$
 $(\text{poincare-between } x \ x1 \ x2 \vee$
 $\text{poincare-between } x1 \ x \ x2 \vee$
 $\text{poincare-between } x1 \ x2 \ x)) \wedge$
 $(\forall a' \in \{z. z \in \text{unit-disc}\}.$
 $a \neq a1 \wedge$
 $a \neq a2 \wedge$
 $a' \neq a \wedge$
 $(\exists x \in \{z. z \in \text{unit-disc}\}.$
 $\text{poincare-between } a1 \ x \ a2 \wedge$
 $x \neq a1 \wedge$
 $x \neq a2 \wedge$
 $(\text{poincare-between } a \ a' \ x \vee$
 $\text{poincare-between } a \ x \ a')) \longrightarrow$
 $(\exists x \in \{z. z \in \text{unit-disc}\}.$
 $(\text{poincare-between } a \ x \ a' \vee$
 $\text{poincare-between } a \ a' \ x) \wedge$
 $(\text{poincare-between } x \ x1 \ x2 \vee$
 $\text{poincare-between } x1 \ x \ x2 \vee$
 $\text{poincare-between } x1 \ x2 \ x)))$
using $**(1, 2)$
by auto
qed
qed

interpretation *PoincareElementaryTarskiHyperbolic: ElementaryTarskiHyperbolic p-congruent p-between*
proof

12. Continuity

fix $\varphi \ \psi$
assume $\exists a. \forall x. \forall y. \varphi \ x \wedge \psi \ y \longrightarrow \text{p-between } a \ x \ y$
thus $\exists b. \forall x. \forall y. \varphi \ x \wedge \psi \ y \longrightarrow \text{p-between } x \ b \ y$

```
    apply transfer
    using continuity
    by auto
qed
end
```

References

- [1] R. Coghetto. Klein-Beltrami Model. Part I. *Formalized Mathematics*, 26(1):21–32, 2018.
- [2] R. Coghetto. Klein-Beltrami Model. Part II. *Formalized Mathematics*, 26(1):33–48, 2018.
- [3] N. Lobatschewsky. *Geometrische Untersuchungen zur Theorie der Parallellinien*, pages 159–223. Springer Vienna, Vienna, 1985.
- [4] T. J. M. Makarios. A Mechanical Verification of the Independence of Tarski’s Euclidean Axiom. Master’s thesis, Victoria University of Wellington, 2012. Master Thesis.
- [5] F. Marić and D. Simić. Formalizing Complex Plane Geometry. *Annals of Mathematics and Artificial Intelligence*, 74(3-4):271–308, 2015.
- [6] F. Mari and D. Simi. Complex geometry. *Archive of Formal Proofs*, Dec. 2019. http://isa-afp.org/entries/Complex_Geometry.html, Formal proof development.
- [7] W. Schwabhäuser, W. Szmielew, and A. Tarski. *Metamathematische Methoden in der Geometrie*. Springer-Verlag, Berlin, 1983.
- [8] H. Schwerdtfeger. *Geometry of Complex Numbers: Circle Geometry, Moebius Transformation, Non-euclidean Geometry*. Courier Corporation, 1979.