

# Pick's Theorem

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## Abstract

We formalize Pick's theorem for finding the area of a simple polygon whose vertices are integral lattice points [1]. We are inspired by John Harrison's formalization of Pick's theorem in HOL Light [2], but tailor our proof approach to avoid a primary challenge point in his formalization, which is proving that any polygon with more than three vertices can be split (in its interior) by a line between some two vertices. Our formalization involves augmenting the existing geometry libraries in various foundational ways (e.g., by adding the definition of a polygon and formalizing some key properties thereof).

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# 1 Misc. Linear Algebra Setup

**lemma** *vec-scaleR-2*:  $(c::real) *_R ((vector [a, b])::real^2) = vector [a * c, b * c]$   
 $\langle proof \rangle$

**definition** *is-int* :: *real*  $\Rightarrow$  *bool* **where**  
 $is\text{-}int x \longleftrightarrow (\exists n::int. x = n)$

**lemma** *is-int-sum*: *is-int*  $x \wedge$  *is-int*  $y \longrightarrow$  *is-int*  $(x + y)$   
 $\langle proof \rangle$

**lemma** *is-int-minus*: *is-int*  $x \wedge$  *is-int*  $y \longrightarrow$  *is-int*  $(x - y)$   
 $\langle proof \rangle$

**lemma** *is-int-mult*: *is-int*  $x \wedge$  *is-int*  $y \longrightarrow$  *is-int*  $(x * y)$   
 $\langle proof \rangle$

**definition** *integral-vec* :: *real^2*  $\Rightarrow$  *bool* **where**  
 $integral\text{-}vec v \longleftrightarrow (is\text{-}int (v\$1) \wedge is\text{-}int (v\$2))$

**lemma** *integral-vec-sum*: *integral-vec*  $v \wedge$  *integral-vec*  $w \longrightarrow$  *integral-vec*  $(v + w)$   
 $\langle proof \rangle$

**lemma** *integral-vec-minus*: *integral-vec*  $v \longrightarrow$  *integral-vec*  $(-v)$   
 $\langle proof \rangle$

**lemma** *real-2-inner*:  
**shows**  $((vector [a, b])::(real^2)) \cdot ((vector [c, d])::(real^2)) = a*c + b*d$   
 $(\text{is } ?v \cdot ?w = a*c + b*d)$   
 $\langle proof \rangle$

**lemma** *integral-vec-2*:  
**fixes**  $a b :: int$   
**assumes**  $v = vector [a, b]$   
**shows** *integral-vec*  $v$   
 $\langle proof \rangle$

**definition** *matrix-inv* :: *real^2^2*  $\Rightarrow$  *real^2^2*  $\Rightarrow$  *bool* **where**  
 $matrix\text{-}inv A A' \longleftrightarrow (A ** A' = mat 1 \wedge A' ** A = mat 1)$

**lemma** *mat-vec-mult-2*:  
**fixes**  $v :: real^2$  **and**  
 $T :: real^2^2$   
**defines**  $x: x \equiv v\$1$  **and**  $y: y \equiv v\$2$  **and**  
 $a: a \equiv T\$1\$1$  **and**  $b: b \equiv T\$1\$2$  **and**  
 $c: c \equiv T\$2\$1$  **and**  $d: d \equiv T\$2\$2$   
**shows**  $(T *v v) = vector [x*a + y*b, x*c + y*d]$   
 $\langle proof \rangle$

```

definition integral-mat :: real22 ⇒ bool where
  integral-mat T ←→ (forall v. integral-vec v → integral-vec (T *v v))

definition integral-mat-surj :: real22 ⇒ bool where
  integral-mat-surj T ←→ (forall v. integral-vec v → exists w. integral-vec w ∧ T *v w = v)

definition integral-mat-bij :: real22 ⇒ bool where
  integral-mat-bij T ←→ integral-mat T ∧ integral-mat-surj T

lemma integral-mat-integral-vec: integral-mat A → integral-vec v → integral-vec (A *v v)
  ⟨proof⟩

lemma integral-mat-int-entries:
  fixes T :: real22
  assumes integral-mat T
  defines a: a ≡ T$1$1 and b: b ≡ T$1$2 and
    c: c ≡ T$2$1 and d: d ≡ T$2$2
  shows is-int a ∧ is-int b ∧ is-int c ∧ is-int d
  ⟨proof⟩

```

## 2 Integral Bijective Matrix Determinant

```

lemma integral-mat-int-det:
  fixes T :: real22
  assumes integral-mat T
  shows is-int (det T)
  ⟨proof⟩

lemma integral-mat-bij-inv:
  fixes T :: real22
  assumes integral-mat-bij T
  obtains Tinv where invertible T ∧ integral-mat-bij Tinv ∧ matrix-inv T Tinv
  ⟨proof⟩

lemma integral-mat-bij-det-pm1:
  fixes T :: real22
  assumes integral-mat-bij T
  shows det T = 1 ∨ det T = -1
  ⟨proof⟩

end
theory Polygon-Jordan-Curve
imports
  HOL-Analysis.Cartesian-Space
  HOL-Analysis.Path-Connected

```

*Poincare-Bendixson.Poincare-Bendixson  
Integral-Matrix*

begin

### 3 Polygon Definitions

**type-synonym**  $R\text{-to-}R2 = (\text{real} \Rightarrow \text{real}^{\wedge 2})$

**definition**  $\text{closed-path} :: R\text{-to-}R2 \Rightarrow \text{bool}$  **where**  
 $\text{closed-path } g \longleftrightarrow \text{path } g \wedge \text{pathstart } g = \text{pathfinish } g$

**definition**  $\text{path-inside} :: R\text{-to-}R2 \Rightarrow (\text{real}^{\wedge 2}) \text{ set}$  **where**  
 $\text{path-inside } g = \text{inside}(\text{path-image } g)$

**definition**  $\text{path-outside} :: R\text{-to-}R2 \Rightarrow (\text{real}^{\wedge 2}) \text{ set}$  **where**  
 $\text{path-outside } g = \text{outside}(\text{path-image } g)$

**fun**  $\text{make-polygonal-path} :: (\text{real}^{\wedge 2}) \text{ list} \Rightarrow R\text{-to-}R2$  **where**  
 $\text{make-polygonal-path } [] = \text{linepath } 0 \ 0$   
 $\mid \text{make-polygonal-path } [a] = \text{linepath } a \ a$   
 $\mid \text{make-polygonal-path } [a,b] = \text{linepath } a \ b$   
 $\mid \text{make-polygonal-path } (a \ # \ b \ # \ xs) = (\text{linepath } a \ b) \ \text{+++} \ \text{make-polygonal-path } (b \ # \ xs)$

**definition**  $\text{polygonal-path} :: R\text{-to-}R2 \Rightarrow \text{bool}$  **where**  
 $\text{polygonal-path } g \longleftrightarrow g \in \text{make-polygonal-path}^{\wedge}\{xs :: (\text{real}^{\wedge 2}) \text{ list}. \ \text{True}\}$

**definition**  $\text{all-integral} :: (\text{real}^{\wedge 2}) \text{ list} \Rightarrow \text{bool}$  **where**  
 $\text{all-integral } l = (\forall x \in \text{set } l. \ \text{integral-vec } x)$

**definition**  $\text{polygon} :: R\text{-to-}R2 \Rightarrow \text{bool}$  **where**  
 $\text{polygon } g \longleftrightarrow \text{polygonal-path } g \wedge \text{simple-path } g \wedge \text{closed-path } g$

**definition**  $\text{integral-polygon} :: R\text{-to-}R2 \Rightarrow \text{bool}$  **where**  
 $\text{integral-polygon } g \longleftrightarrow$   
 $(\text{polygon } g \wedge (\exists vts. \ g = \text{make-polygonal-path } vts \wedge \text{all-integral } vts))$

**definition**  $\text{make-triangle} :: \text{real}^{\wedge 2} \Rightarrow \text{real}^{\wedge 2} \Rightarrow \text{real}^{\wedge 2} \Rightarrow R\text{-to-}R2$  **where**  
 $\text{make-triangle } a \ b \ c = \text{make-polygonal-path } [a, b, c, a]$

**definition**  $\text{polygon-of} :: R\text{-to-}R2 \Rightarrow (\text{real}^{\wedge 2}) \text{ list} \Rightarrow \text{bool}$  **where**  
 $\text{polygon-of } p \ vts \longleftrightarrow \text{polygon } p \wedge p = \text{make-polygonal-path } vts$

**definition**  $\text{good-linepath} :: \text{real}^{\wedge 2} \Rightarrow \text{real}^{\wedge 2} \Rightarrow (\text{real}^{\wedge 2}) \text{ list} \Rightarrow \text{bool}$  **where**  
 $\text{good-linepath } a \ b \ vts \longleftrightarrow (\text{let } p = \text{make-polygonal-path } vts \text{ in}$   
 $a \neq b \wedge \{a, b\} \subseteq \text{set } vts \wedge \text{path-inside } (\text{linepath } a \ b) \subseteq \text{path-inside } p \cup \{a, b\})$

**definition**  $\text{good-polygonal-path} :: \text{real}^{\wedge 2} \Rightarrow (\text{real}^{\wedge 2}) \text{ list} \Rightarrow \text{real}^{\wedge 2} \Rightarrow (\text{real}^{\wedge 2}) \text{ list}$

```

 $\Rightarrow \text{bool where}$ 
 $\text{good-polygonal-path } a \text{ cutvts } b \text{ vts} \longleftrightarrow ($ 
 $\quad \text{let } p = \text{make-polygonal-path } \text{vts} \text{ in}$ 
 $\quad \text{let } p\text{-cut} = \text{make-polygonal-path } ([a] @ \text{cutvts} @ [b]) \text{ in}$ 
 $\quad (a \neq b \wedge \{a, b\} \subseteq \text{set vts} \wedge \text{path-image } (p\text{-cut}) \subseteq \text{path-inside } p \cup \{a, b\} \wedge$ 
 $\quad \text{loop-free } p\text{-cut}))$ 

```

## 4 Jordan Curve Theorem for Polygons

```

definition inside-outside :: R-to-R2  $\Rightarrow (\text{real}^2) \text{ set} \Rightarrow (\text{real}^2) \text{ set} \Rightarrow \text{bool where}$ 
inside-outside p ins outs  $\longleftrightarrow$ 
 $(\text{ins} \neq \{\} \wedge \text{open ins} \wedge \text{connected ins} \wedge$ 
 $\text{outs} \neq \{\} \wedge \text{open outs} \wedge \text{connected outs} \wedge$ 
 $\text{bounded ins} \wedge \neg \text{bounded outs} \wedge$ 
 $\text{ins} \cap \text{outs} = \{\} \wedge \text{ins} \cup \text{outs} = - \text{path-image } p \wedge$ 
 $\text{frontier ins} = \text{path-image } p \wedge \text{frontier outs} = \text{path-image } p)$ 

```

**lemma** *Jordan-inside-outside-real2*:

```

fixes p :: real  $\Rightarrow \text{real}^2$ 
assumes simple-path p pathfinish p = pathstart p
shows inside(path-image p)  $\neq \{\}$   $\wedge$ 
 $\text{open}(\text{inside}(\text{path-image } p)) \wedge$ 
 $\text{connected}(\text{inside}(\text{path-image } p)) \wedge$ 
 $\text{outside}(\text{path-image } p) \neq \{\} \wedge$ 
 $\text{open}(\text{outside}(\text{path-image } p)) \wedge$ 
 $\text{connected}(\text{outside}(\text{path-image } p)) \wedge$ 
 $\text{bounded}(\text{inside}(\text{path-image } p)) \wedge$ 
 $\neg \text{bounded}(\text{outside}(\text{path-image } p)) \wedge$ 
 $\text{inside}(\text{path-image } p) \cap \text{outside}(\text{path-image } p) = \{\} \wedge$ 
 $\text{inside}(\text{path-image } p) \cup \text{outside}(\text{path-image } p) =$ 
 $- \text{path-image } p \wedge$ 
 $\text{frontier}(\text{inside}(\text{path-image } p)) = \text{path-image } p \wedge$ 
 $\text{frontier}(\text{outside}(\text{path-image } p)) = \text{path-image } p$ 

```

*{proof}*

**lemma** *inside-outside-polygon*:

```

fixes p :: R-to-R2
assumes polygon: polygon p
shows inside-outside p (path-inside p) (path-outside p)

```

*{proof}*

**lemma** *inside-outside-unique*:

```

fixes p :: R-to-R2
assumes polygon p
assumes io1: inside-outside p inside1 outside1
assumes io2: inside-outside p inside2 outside2
shows inside1 = inside2  $\wedge$  outside1 = outside2

```

$\langle proof \rangle$

**lemma** *polygon-jordan-curve*:  
  **fixes**  $p :: R\text{-to-}R^2$   
  **assumes** *polygon p*  
  **obtains** *inside outside* **where**  
    *inside-outside p inside outside*  
 $\langle proof \rangle$

**lemma** *connected-component-image*:  
  **fixes**  $f :: 'a\text{::euclidean-space} \Rightarrow 'b\text{::euclidean-space}$   
  **assumes** *linear f bij f*  
  **shows**  $f`(\text{connected-component-set } S x) = \text{connected-component-set } (f`S) (f x)$   
 $\langle proof \rangle$

**lemma** *bounded-map*:  
  **fixes**  $f :: 'a\text{::euclidean-space} \Rightarrow 'b\text{::euclidean-space}$   
  **assumes** *linear f bij f*  
  **shows** *bounded (f`S) = bounded S*  
 $\langle proof \rangle$

**lemma** *inside-bijective-linear-image*:  
  **fixes**  $f :: 'a\text{::euclidean-space} \Rightarrow 'b\text{::euclidean-space}$   
  **fixes**  $c :: \text{real} \Rightarrow 'a$   
  **assumes** *c-simple:path c*  
  **assumes** *linear f bij f*  
  **shows** *inside (f`(\text{path-image } c)) = f`(\text{inside(path-image } c))*  
 $\langle proof \rangle$

**lemma** *bij-image-intersection*:  
  **assumes** *path-image c1 ∩ path-image c2 = S*  
  **assumes** *bij f*  
  **assumes**  $c \in \text{path-image } (f \circ c1) \cap \text{path-image } (f \circ c2)$   
  **shows**  $c \in f`S$   
 $\langle proof \rangle$

**theorem (in c1-on-open-R2) split-inside-simple-closed-curve-locale**:  
  **fixes**  $c :: \text{real} \Rightarrow 'a$   
  **assumes** *c1-simple:simple-path c1 and c1-start: pathstart c1 = a and c1-end: pathfinish c1 = b*  
  **assumes** *c2-simple: simple-path c2 and c2-start: pathstart c2 = a and c2-end: pathfinish c2 = b*  
  **assumes** *c-simple: simple-path c and c-start: pathstart c = a and c-end: pathfinish c = b*  
  **assumes** *a-neq-b: a ≠ b*  
    **and** *c1c2: path-image c1 ∩ path-image c2 = {a,b}*  
    **and** *c1c: path-image c1 ∩ path-image c = {a,b}*

```

and c2c: path-image c2 ∩ path-image c = {a,b}
and ne-12: path-image c ∩ inside(path-image c1 ∪ path-image c2) ≠ {}
obtains inside(path-image c1 ∪ path-image c) ∩ inside(path-image c2 ∪ path-image
c) = {}
           inside(path-image c1 ∪ path-image c) ∪ inside(path-image c2 ∪ path-image
c) ∪
           (path-image c - {a,b}) = inside(path-image c1 ∪ path-image c2)
⟨proof⟩

lemma split-inside-simple-closed-curve-real2:
fixes c :: real ⇒ real^2
assumes c1-simple:simple-path c1 and c1-start: pathstart c1 = a and c1-end:
pathfinish c1 = b
assumes c2-simple: simple-path c2 and c2-start: pathstart c2 = a and c2-end:
pathfinish c2 = b
assumes c-simple: simple-path c and c-start: pathstart c = a and c-end: pathfin-
ish c = b
assumes a-neq-b: a ≠ b
and c1c2: path-image c1 ∩ path-image c2 = {a,b}
and c1c: path-image c1 ∩ path-image c = {a,b}
and c2c: path-image c2 ∩ path-image c = {a,b}
and ne-12: path-image c ∩ inside(path-image c1 ∪ path-image c2) ≠ {}
obtains inside(path-image c1 ∪ path-image c) ∩ inside(path-image c2 ∪ path-image
c) = {}
           inside(path-image c1 ∪ path-image c) ∪ inside(path-image c2 ∪ path-image
c) ∪
           (path-image c - {a,b}) = inside(path-image c1 ∪ path-image c2)
⟨proof⟩

end
theory Polygon-Lemmas
imports
  Polygon-Jordan-Curve
  HOL-Library.Sublist
  HOL.Set-Interval
  HOL.Fun

begin

```

## 5 Properties of make\_polygonal\_path, pathstart and pathfinish of a polygon

```

lemma make-polygonal-path-induct[case-names Empty Single Two Multiple]:
fixes ell :: (real^2) list
assumes empty: ⋀ell. ell = [] ⇒ P ell
and single: ⋀ell. [length ell = 1] ⇒ P ell
and two: ⋀ell. [length ell = 2] ⇒ P ell
and multiple: ⋀ell.

```

```

 $\llbracket \text{length } ell > 2;$ 
 $P ((ell!0), (ell!1))$ ;
 $P ((ell!1) \# (drop 2 ell)) \rrbracket \implies P ell$ 
shows  $P ell$ 
 $\langle proof \rangle$ 

lemma make-polygonal-path-gives-path:
fixes  $v :: (\text{real}^2)^{\text{list}}$ 
shows path (make-polygonal-path  $v$ )
 $\langle proof \rangle$ 

corollary polygonal-path-is-path:
fixes  $g :: R\text{-to-}R2$ 
assumes polygonal-path  $g$ 
shows path  $g$ 
 $\langle proof \rangle$ 

lemma polygon-to-polygonal-path:
fixes  $p :: R\text{-to-}R2$ 
assumes polygon  $p$ 
obtains  $ell$  where  $p = \text{make-polygonal-path } ell$ 
 $\langle proof \rangle$ 

lemma polygon-pathstart:
fixes  $g :: R\text{-to-}R2$ 
assumes  $l \neq []$ 
assumes  $g = \text{make-polygonal-path } l$ 
shows pathstart  $g = l!0$ 
 $\langle proof \rangle$ 

lemma polygon-pathfinish:
fixes  $g :: R\text{-to-}R2$ 
assumes  $l \neq []$ 
assumes  $g = \text{make-polygonal-path } l$ 
shows pathfinish  $g = l!(\text{length } l - 1)$ 
 $\langle proof \rangle$ 

lemma make-polygonal-path-image-property:
assumes length  $vts \geq 2$ 
assumes p-is-path:  $x \in \text{path-image}(\text{make-polygonal-path } vts)$ 
shows  $\exists k < \text{length } vts - 1. x \in \text{path-image}(\text{linepath}(vts ! k) (vts ! (k + 1)))$ 
 $\langle proof \rangle$ 

lemma linepaths-subset-make-polygonal-path-image:
assumes length  $vts \geq 2$ 
assumes  $k < \text{length } vts - 1$ 
shows path-image (linepath ( $vts ! k$ ) ( $vts ! (k + 1)$ ))  $\subseteq \text{path-image}(\text{make-polygonal-path } vts)$ 

```

$\langle proof \rangle$

**lemma** *vertices-on-path-image*: **shows** set *vts*  $\subseteq$  *path-image* (*make-polygonal-path* *vts*)  
 $\langle proof \rangle$

**lemma** *path-image-cons-union*:  
**assumes** *p* = *make-polygonal-path* *vts*  
**assumes** *p'* = *make-polygonal-path* *vts'*  
**assumes** *vts' ≠ []*  
**assumes** *vts* = *a* # *vts'*  $\wedge$  *b* = *vts'!0*  
**shows** *path-image p* = *path-image* (*linepath a b*)  $\cup$  *path-image p'*  
 $\langle proof \rangle$

**lemma** *polygonal-path-image-linepath-union*:  
**assumes** *p* = *make-polygonal-path* *vts*  
**assumes** *n* = *length vts*  
**assumes** *n ≥ 2*  
**shows** *path-image p* = ( $\bigcup$  {*path-image* (*linepath* (*vts!i*) (*vts!(i+1)*)) | *i*. *i ≤ n - 2*})  
 $\langle proof \rangle$

## 6 Loop Free Properties

**lemma** *constant-linepath-is-not-loop-free*:  
**shows**  $\neg(\text{loop-free } ((\text{linepath } a \ a)::\text{real} \Rightarrow \text{real}^{\wedge 2}))$   
 $\langle proof \rangle$

**lemma** *doubling-back-is-not-loop-free*:  
**assumes** *a ≠ b*  
**shows**  $\neg(\text{loop-free } ((\text{make-polygonal-path } [a, b, a])::\text{real} \Rightarrow \text{real}^{\wedge 2}))$   
 $\langle proof \rangle$

**lemma** *not-loop-free-first-component*:  
**assumes**  $\neg(\text{loop-free } p1)$   
**shows**  $\neg(\text{loop-free } (p1+++p2))$   
 $\langle proof \rangle$

**lemma** *not-loop-free-second-component*:  
**assumes** *pathfinish-pathstart*: *pathfinish p1* = *pathstart p2*  
**assumes**  $\neg(\text{loop-free } p2)$   
**shows**  $\neg(\text{loop-free } (p1+++p2))$   
 $\langle proof \rangle$

**lemma** *loop-free-subpath*:  
**assumes** *path p*  
**assumes** *u-and-v*: *u ∈ {0..1}* *v ∈ {0..1}* *u < v*  
**assumes**  $\neg(\text{loop-free } (\text{subpath } u \ v \ p))$   
**shows**  $\neg(\text{loop-free } p)$

$\langle proof \rangle$

**lemma** *loop-free-associative*:

**assumes** *path p*

**assumes** *path q*

**assumes** *path r*

**assumes** *pathfinish p = pathstart q*

**assumes** *pathfinish q = pathstart r*

**shows**  $\neg (\text{loop-free } ((p \text{ +++ } q) \text{ +++ } r)) \longleftrightarrow \neg (\text{loop-free } (p \text{ +++ } (q \text{ +++ } r)))$

$\langle proof \rangle$

**lemma** *polygon-at-least-3-vertices*:

**assumes** *polygon p and*

*p = make-polygonal-path vts*

**shows** *card (set vts) ≥ 3*

$\langle proof \rangle$

**lemma** *polygon-vertices-length-at-least-4*:

**assumes** *polygon p and*

*p = make-polygonal-path vts*

**shows** *length vts ≥ 4*

$\langle proof \rangle$

**lemma** *linepath-loop-free*:

**assumes** *a ≠ b*

**shows** *loop-free (linepath a b)*

$\langle proof \rangle$

## 7 Explicit Linepath Characterization of Polygonal Paths

**lemma** *triangle-linepath-images*:

**fixes** *x :: real*

**assumes** *vts = [a, b, c]*

**assumes** *p = make-polygonal-path vts*

**shows** *x ∈ {0..1/2} ⇒ p x = ((linepath a b)) (2\*x)*

*x ∈ {1/2..1} ⇒ p x = ((linepath b c)) (2\*x - 1)*

$\langle proof \rangle$

**lemma** *polygon-linepath-images1*:

**fixes** *n:: nat*

**assumes** *n ≥ 3*

**assumes** *length ell = n*

**assumes** *x ∈ {0..1/2}*

**shows** *make-polygonal-path ell x = ((linepath (ell ! 0) (ell ! 1))) (2\*x)*

$\langle proof \rangle$

```

lemma sum-insert [simp]:
  assumes  $x \notin F$  and finite  $F$ 
  shows  $(\sum y \in \text{insert } x F. P y) = (\sum y \in F. P y) + P x$ 
   $\langle proof \rangle$ 

lemma sum-of-index-diff [simp]:
  fixes  $f :: nat \Rightarrow 'a :: \text{comm-monoid-add}$ 
  shows  $(\sum i \in \{a..<a+b\}. f(i-a)) = (\sum i \in \{..<b\}. f(i))$ 
   $\langle proof \rangle$ 

lemma sum-of-index-diff2 [simp]:
  fixes  $f :: nat \Rightarrow 'a :: \text{comm-monoid-add}$ 
  shows  $(\sum i \in \{a+c..b+c\}. f(i)) = (\sum i \in \{a..b\}. f(i+c))$ 
   $\langle proof \rangle$ 

lemma sum-split [simp]:
  fixes  $f :: nat \Rightarrow 'a :: \text{comm-monoid-add}$ 
  assumes  $c \in \{a..b\}$ 
  shows  $(\sum i \in \{a..b\}. f i) = (\sum i \in \{a..c\}. f i) + (\sum i \in \{c+1..b\}. f i)$ 
   $\langle proof \rangle$ 

lemma summation-helper:
  fixes  $x :: real$ 
  fixes  $k :: nat$ 
  assumes  $1 \leq k$ 
  shows  $(2 :: real) * (\sum i = 1..k. 1 / 2^i) - 1 = (\sum i = 1..(k-1). (1 / (2^i)))$ 
   $\langle proof \rangle$ 

lemma polygon-linepath-images2:
  fixes  $n k :: nat$ 
  fixes  $ell :: (real^2) list$ 
  fixes  $f :: nat \Rightarrow real \Rightarrow real$ 
  assumes  $n \geq 3$ 
  assumes  $0 \leq k \wedge k \leq n - 3$ 
  assumes  $\text{length } ell = n$ 
  assumes  $p :: p = \text{make-polygonal-path } ell$ 
  assumes  $f = (\lambda k x. (x - (\sum i \in \{1..k\}. 1 / (2^i))) * (2^{k+1}))$ 
  assumes  $x \in \{(\sum i \in \{1..k\}. 1 / (2^i))..(\sum i \in \{1..(k+1)\}. 1 / (2^i))\}$ 
  shows  $p x = ((\text{linepath} (ell ! k) (ell ! (k+1)) (f k x)))$ 
   $\langle proof \rangle$ 

lemma polygon-linepath-images3:
  fixes  $n k :: nat$ 
  fixes  $ell :: (real^2) list$ 
  assumes  $n \geq 3$ 
  assumes  $\text{length } ell = n$ 
  assumes  $p = \text{make-polygonal-path } ell$ 
  assumes  $x \in \{(\sum i \in \{1..n-2\}. 1 / (2^i))..1\}$ 

```

```

assumes  $f = (\lambda x. (x - (\sum i \in \{1..n-2\}. 1/(2^i))) * (2^{n-2}))$ 
shows  $p x = (\text{linepath} (\text{ell} ! (n-2)) (\text{ell} ! (n-1))) (f x)$ 
⟨proof⟩

```

## 8 A Triangle is a Polygon

```

lemma not-collinear-linepaths-intersect-helper:
assumes not-collinear:  $\neg \text{collinear } \{a,b,c\}$ 
assumes  $0 \leq k1$ 
assumes  $k1 \leq 1$ 
assumes  $0 \leq k2$ 
assumes  $k2 \leq 1$ 
assumes eo:  $k2 = 0 \implies k1 \neq 1$ 
shows  $\neg ((\text{linepath } a b) k1 = (\text{linepath } b c) k2)$ 
⟨proof⟩

```

```

lemma not-collinear-linepaths-intersect-helper-2:
assumes not-collinear:  $\neg \text{collinear } \{a,b,c\}$ 
assumes  $0 \leq k1$ 
assumes  $k1 \leq 1$ 
assumes  $0 \leq k2$ 
assumes  $k2 \leq 1$ 
assumes eo:  $k1 = 0 \implies k2 \neq 1$ 
shows  $\neg ((\text{linepath } a b) k1 = (\text{linepath } c a) k2)$ 
⟨proof⟩

```

```

lemma not-collinear-loopfree-path:  $\bigwedge a b c::\text{real}^2. \neg \text{collinear } \{a,b,c\} \implies \text{loop-free} ((\text{linepath } a b) +\!\!+\!+ (\text{linepath } b c))$ 
⟨proof⟩

```

```

lemma triangle-is-polygon:  $\bigwedge a b c. \neg \text{collinear } \{a,b,c\} \implies \text{polygon } (\text{make-triangle } a b c)$ 
⟨proof⟩

```

```

lemma have-wraparound-vertex:
assumes polygon  $p$ 
assumes  $p = \text{make-polygonal-path } vts$ 
shows  $vts = (\text{take} (\text{length } vts - 1) vts) @ [vts ! 0]$ 
⟨proof⟩

```

```

lemma polygon-at-least-3-vertices-wraparound:
assumes polygon  $p$ 
assumes  $p = \text{make-polygonal-path } vts$ 
shows  $\text{card} (\text{set} (\text{take} (\text{length } vts - 1) vts)) \geq 3$ 
⟨proof⟩

```

## 9 Polygon Vertex Rotation

```

definition rotate-polygon-vertices:: 'a list  $\Rightarrow$  nat  $\Rightarrow$  'a list
  where rotate-polygon-vertices ell i =
    (let ell1 = rotate i (butlast ell) in ell1 @ [ell1 ! 0])

lemma rotate-polygon-vertices-same-set:
  assumes polygon (make-polygonal-path vts)
  shows set (rotate-polygon-vertices vts i) = set vts
  ⟨proof⟩

lemma arb-rotation-as-single-rotation:
  fixes i:: nat
  shows rotate-polygon-vertices vts (Suc i) = rotate-polygon-vertices (rotate-polygon-vertices vts i) 1
  ⟨proof⟩

lemma rotation-sum:
  fixes i j :: nat
  shows rotate-polygon-vertices vts (i + j) = rotate-polygon-vertices (rotate-polygon-vertices vts i) j
  ⟨proof⟩

lemma rotated-polygon-vertices-helper:
  fixes p :: R-to-R $\circlearrowleft$ 
  assumes poly-p: polygon p
  assumes p-is-path: p = make-polygonal-path vts
  assumes p'-is: p' = make-polygonal-path (rotate-polygon-vertices vts 1)
  shows (vts ! 0) = (rotate-polygon-vertices vts 1) ! (length (rotate-polygon-vertices vts 1) - 2)
    (rotate-polygon-vertices vts 1) ! (length (rotate-polygon-vertices vts 1) - 1)
  = (vts ! 1)
  ⟨proof⟩

lemma rotate-polygon-vertices-same-length:
  fixes vts :: (real $\circlearrowleft$ ) list
  assumes length vts  $\geq$  1
  shows length vts = length (rotate-polygon-vertices vts i)
  ⟨proof⟩

lemma rotated-polygon-vertices-helper2:
  assumes len-gteq: length vts  $\geq$  2
  assumes i < length vts - 1
  assumes hd vts = last vts
  shows (rotate-polygon-vertices vts 1) ! i = vts ! (i+1)
  ⟨proof⟩

lemma polygon-rotation-t-translation1:
  assumes polygon-of p vts

```

```

assumes  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$ 
      (is  $p' = \text{make-polygonal-path} ?vts'$ )
assumes  $x' \in \{(\sum i \in \{1..k\}. 1/(2^i))..(\sum i \in \{1..k+1\}. 1/(2^i))\}$ 
assumes  $n = \text{length } vts$ 
assumes  $0 \leq k \wedge k \leq n - 4$ 
assumes  $l = x' - (\sum i \in \{1..k\}. 1/(2^i))$ 
assumes  $x = l/2 + (\sum i \in \{1..(k+1)\}. 1/(2^i))$ 
shows  $x \in \{(\sum i \in \{1..k+1\}. 1/(2^i))..(\sum i \in \{1..k+2\}. 1/(2^i))\}$ 

$$p' x' = p x$$

⟨proof⟩

```

```

lemma polygon-rotation-t-translation1-strict:
assumes polygon-of  $p$  vts
assumes  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$ 
      (is  $p' = \text{make-polygonal-path} ?vts'$ )
assumes  $x' \in \{(\sum i \in \{1..k\}. 1/(2^i))..<(\sum i \in \{1..k+1\}. 1/(2^i))\}$ 
assumes  $n = \text{length } vts$ 
assumes  $0 \leq k \wedge k \leq n - 4$ 
assumes  $l = x' - (\sum i \in \{1..k\}. 1/(2^i))$ 
assumes  $x = l/2 + (\sum i \in \{1..(k+1)\}. 1/(2^i))$ 
shows  $x \in \{(\sum i \in \{1..k+1\}. 1/(2^i))..<(\sum i \in \{1..k+2\}. 1/(2^i))\}$ 

$$p' x' = p x$$

⟨proof⟩

```

```

lemma polygon-rotation-t-translation2:
assumes polygon-of  $p$  vts
assumes  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$ 
      (is  $p' = \text{make-polygonal-path} ?vts'$ )
assumes  $n = \text{length } vts$ 
assumes  $x' \in \{(\sum i \in \{1..(n-3)\}. 1/(2^i))..(\sum i \in \{1..(n-2)\}. 1/(2^i))\}$ 
assumes  $x = x' + 1/(2^{(n-2)})$ 
shows  $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$ 

$$p' x' = p x$$

⟨proof⟩

```

```

lemma polygon-rotation-t-translation2-strict:
assumes polygon-of  $p$  vts
assumes  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$ 
      (is  $p' = \text{make-polygonal-path} ?vts'$ )
assumes  $n = \text{length } vts$ 
assumes  $x' \in \{(\sum i \in \{1..(n-3)\}. 1/(2^i))..<(\sum i \in \{1..(n-2)\}. 1/(2^i))\}$ 
assumes  $x = x' + 1/(2^{(n-2)})$ 
shows  $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..<1\}$ 

$$p' x' = p x$$

⟨proof⟩

```

```

lemma polygon-rotation-t-translation3:
assumes polygon-of  $p$  vts

```

```

assumes  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$ 
  (is  $p' = \text{make-polygonal-path} ?vts'$ )
assumes  $x' \in \{(\sum i \in \{1..n-2\}. \ 1/(2^i))..1\}$ 
assumes  $n = \text{length } vts$ 
assumes  $l = x' - (\sum i \in \{1..n-2\}. \ 1/(2^i))$ 
assumes  $x = l * (2^{(n-3)})$ 
shows  $x \in \{0..1/2\}$ 

$$p' x' = p x$$

⟨proof⟩

```

```

lemma polygon-rotation-t-translation3-strict:
assumes polygon-of  $p \ vts$ 
assumes  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$ 
  (is  $p' = \text{make-polygonal-path} ?vts'$ )
assumes  $x' \in \{(\sum i \in \{1..n-2\}. \ 1/(2^i))..<1\}$ 
assumes  $n = \text{length } vts$ 
assumes  $l = x' - (\sum i \in \{1..n-2\}. \ 1/(2^i))$ 
assumes  $x = l * (2^{(n-3)})$ 
shows  $x \in \{0..<1/2\}$ 

$$p' x' = p x$$

⟨proof⟩

```

```

lemma f-gteq-0-sum-gt:  $\bigwedge f::nat \Rightarrow \text{real}. \ (\bigwedge i::nat. \ (f i) > 0) \implies a > b \implies (\sum i = 1..a. \ (f i)) > (\sum i = 1..b. \ (f i))$  for  $a \ b :: nat$ 
⟨proof⟩

```

```

lemma rotation-intervals-disjoint:
assumes  $k1 \neq k2$ 
shows  $\{\sum i = 1..k1. \ 1 / (2^i)::real\}..<\sum i = 1..k1+1. \ 1 / (2^i)\} \cap \{\sum i = 1..k2. \ 1 / (2^i)::real\}..<\sum i = 1..k2+1. \ 1 / (2^i)\} = \{\}$ 
⟨proof⟩

```

```

lemma bounding-interval-helper1:
shows  $(\sum i = 1..k. \ 1 / (2^i)::real) = (2^k - 1)/(2^k)$ 
⟨proof⟩

```

```

lemma bounding-interval-helper2:
fixes  $x :: \text{real}$ 
assumes  $x \in \{0..<1\}$ 
shows  $\exists k. \ x < (\sum i = 1..k. \ 1 / (2^i)::real)$ 
⟨proof⟩

```

```

lemma bounding-interval-for-reals-btw01:
fixes  $x::\text{real}$ 
assumes  $x \in \{0..<1\}$ 
shows  $\exists k. \ x \in \{(\sum i \in \{1..k\}. \ 1/(2^i)::real)\}..<(\sum i \in \{1..(k+1)\}. \ 1/(2^i))\}$ 
⟨proof⟩

```

```

lemma all-rotation-intervals-between-0and1:

```

**shows**  $\{(\sum i \in \{1..k\}. 1/(2^i :: \text{real})) .. (\sum i \in \{1..(k+1)\}. 1/(2^i))\} \subseteq \{0..<1\}$   
 $\langle \text{proof} \rangle$

**lemma** *all-rotation-intervals-between-0and1-strict*:  
**shows**  $\{(\sum i \in \{1..k\}. 1/(2^i :: \text{real})) .. < (\sum i \in \{1..(k+1)\}. 1/(2^i))\} \subseteq \{0..<1\}$   
 $\langle \text{proof} \rangle$

**lemma** *one-polygon-rotation-is-loop-free*:  
**assumes** *polygon-of p vts*  
**assumes**  $p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$   
**(is**  $p' = \text{make-polygonal-path}(\text{?vts}')$   
**shows** *loop-free p'*  
 $\langle \text{proof} \rangle$

**lemma** *one-rotation-is-polygon*:  
**fixes**  $p :: R\text{-to-}R^2$   
**fixes**  $i :: \text{nat}$   
**assumes**  $\text{poly-}p: \text{polygon } p \text{ and}$   
 $p\text{-is-path}: p = \text{make-polygonal-path } vts \text{ and}$   
 $p'\text{-is}: p' = \text{make-polygonal-path}(\text{rotate-polygon-vertices } vts \ 1)$   
**(is**  $p' = \text{make-polygonal-path}(\text{?vts}')$   
**shows** *polygon p'*  
 $\langle \text{proof} \rangle$

**lemma** *rotation-is-polygon*:  
**fixes**  $p :: R\text{-to-}R^2$   
**fixes**  $i :: \text{nat}$   
**assumes** *polygon p and*  
 $p = \text{make-polygonal-path } vts$   
**shows** *polygon (make-polygonal-path (rotate-polygon-vertices vts i))*  
 $\langle \text{proof} \rangle$

**lemma** *polygon-rotate-mod*:  
**fixes**  $vts :: (\text{real}^2) \text{ list}$   
**assumes**  $n = \text{length } vts$   
**assumes**  $n \geq 2$   
**assumes**  $\text{hd } vts = \text{last } vts$   
**shows**  $\text{rotate-polygon-vertices } vts (n - 1) = vts$   
 $\langle \text{proof} \rangle$

**lemma** *polygon-rotate-mod-arg*:  
**fixes**  $vts :: (\text{real}^2) \text{ list}$   
**assumes**  $n = \text{length } vts$   
**assumes**  $n \geq 2$   
**assumes**  $\text{hd } vts = \text{last } vts$   
**shows**  $\text{rotate-polygon-vertices } vts ((n - 1) * i) = vts$   
 $\langle \text{proof} \rangle$

**lemma** *unrotation-is-polygon*:

```

fixes p :: R-to-R2
fixes i:: nat
assumes polygon (make-polygonal-path (rotate-polygon-vertices vts i))
    (is polygon (make-polygonal-path ?vts'))
    p = make-polygonal-path vts
    hd vts = last vts
shows polygon p
⟨proof⟩

lemma rotated-polygon-vertices:
assumes vts' = rotate-polygon-vertices vts j
assumes hd vts = last vts
assumes length vts ≥ 2
assumes j ≤ i ∧ i < length vts
shows vts ! i = vts' ! (i - j)
⟨proof⟩

lemma polygon-path-image:
assumes poly-p: polygon p
assumes p-is-path: p = make-polygonal-path vts
shows path-image p = p` {0 ..< 1}
⟨proof⟩

lemma polygon-vts-one-rotation:
fixes p :: R-to-R2
assumes poly-p: polygon p and
    p-is-path: p = make-polygonal-path vts and
    p'-is: p' = make-polygonal-path (rotate-polygon-vertices vts 1)
shows path-image p = path-image p'
⟨proof⟩

lemma polygon-vts-arb-rotation:
fixes p :: R-to-R2
assumes polygon p and
    p = make-polygonal-path vts
shows path-image p = path-image (make-polygonal-path (rotate-polygon-vertices
vts i))
⟨proof⟩

```

## 10 Translating a Polygon

```

lemma linopath-translation:
    linopath ((λx. x + u) a) ((λx. x + u) b) = (λx. x + u) ∘ (linopath a b)
⟨proof⟩

lemma make-polygonal-path-translate:
assumes length vts ≥ 2
shows make-polygonal-path (map (λx. x + u) vts) = (λx. x + u) ∘ (make-polygonal-path
vts)

```

$\langle proof \rangle$

**lemma** *translation-is-polygon*:  
  **assumes** *polygon-of p vts*  
  **shows** *polygon-of ((λx. x + u) ∘ p) (map (λx. x + u) vts)* (**is polygon-of ?p'**  
  *?vts'*)  
 $\langle proof \rangle$

## 11 Misc. properties

**lemma** *tail-of-loop-free-polygonal-path-is-loop-free*:  
  **assumes** *loop-free (make-polygonal-path (x#tail))* (**is loop-free ?p**) **and**  
    *length tail ≥ 2*  
  **shows** *loop-free (make-polygonal-path tail)* (**is loop-free ?p'**)  
 $\langle proof \rangle$

**lemma** *tail-of-simple-polygonal-path-is-simple*:  
  **assumes** *simple-path (make-polygonal-path (x#tail))* (**is simple-path ?p**) **and**  
    *length tail ≥ 2*  
  **shows** *simple-path (make-polygonal-path tail)* (**is simple-path ?p'**)  
 $\langle proof \rangle$

**lemma** *interior-vtx-in-path-image-interior*:  
  **fixes** *vts :: (real^2) list*  
  **assumes** *x ∈ set (butlast (drop 1 vts))*  
  **shows**  $\exists t. t \in \{0 <.. < 1\} \wedge (\text{make-polygonal-path vts}) t = x$   
 $\langle proof \rangle$

**lemma** *loop-free-polygonal-path-vts-distinct*:  
  **assumes** *loop-free (make-polygonal-path vts)*  
  **shows** *distinct (butlast vts)*  
 $\langle proof \rangle$

**lemma** *loop-free-polygonal-path-vts-drop1-distinct*:  
  **assumes** *loop-free (make-polygonal-path vts)*  
  **shows** *distinct (drop 1 vts)*  
 $\langle proof \rangle$

**lemma** *simple-polygonal-path-vts-distinct*:  
  **assumes** *simple-path (make-polygonal-path vts)*  
  **shows** *distinct (butlast vts)*  
 $\langle proof \rangle$

**lemma** *edge-subset-path-image*:  
  **assumes** *p = make-polygonal-path vts and*  
    *(i::int) ∈ {0..<((length vts) - 1)}* **and**  
    *x = vts!i and*

$y = vts!(i+1)$   
**shows**  $\text{path-image}(\text{linepath } x \ y) \subseteq \text{path-image } p$  (**is**  $?xy-img \subseteq ?p-img$ )  
 $\langle proof \rangle$

## 12 Properties of Sublists of Polygonal Path Vertex Lists

```

lemma make-polygonal-path-image-append-var:
  assumes length vts1  $\geq 2$ 
  shows path-image (make-polygonal-path (vts1 @ [v])) = path-image (make-polygonal-path
  vts1 +++ (linepath (vts1 ! (length vts1 - 1)) v))
   $\langle proof \rangle$ 

lemma make-polygonal-path-image-append-helper:
  assumes length vts1  $\geq 1 \wedge$  length vts2  $\geq 1$ 
  shows path-image (make-polygonal-path (vts1 @ [v] @ [v] @ vts2)) = path-image
  (make-polygonal-path (vts1 @ [v] @ vts2))
   $\langle proof \rangle$ 

lemma make-polygonal-path-image-append:
  assumes length vts1  $\geq 2 \wedge$  length vts2  $\geq 2$ 
  shows path-image (make-polygonal-path (vts1 @ vts2)) = path-image (make-polygonal-path
  vts1 +++ (linepath (vts1 ! (length vts1 - 1)) (vts2 ! 0)) +++ make-polygonal-path
  vts2)
   $\langle proof \rangle$ 

lemma make-polygonal-path-image-append-alt:
  assumes p = make-polygonal-path vts
  assumes p1 = make-polygonal-path vts1
  assumes p2 = make-polygonal-path vts2
  assumes last vts1 = hd vts2
  assumes length vts1  $\geq 2 \wedge$  length vts2  $\geq 2$ 
  assumes vts = vts1 @ (tl vts2)
  shows path-image p = path-image (p1 +++ p2)
   $\langle proof \rangle$ 

lemma cont-incr-interval-image:
  fixes f :: real  $\Rightarrow$  real
  assumes a  $\leq$  b
  assumes continuous-on {a..b} f
  assumes  $\forall x \in \{a..b\}. \forall y \in \{a..b\}. x \leq y \longrightarrow f x \leq f y$ 
  shows f'{a..b} = {f a..f b}
   $\langle proof \rangle$ 

lemma two-x-minus-one-image:
  assumes f = ( $\lambda x::real. 2*x - 1$ )
  assumes a  $\leq$  b
  shows f'{a..b} = {f a..f b}

```

$\langle proof \rangle$

```
lemma vts-split-path-image:  
  assumes p = make-polygonal-path vts  
  assumes p1 = make-polygonal-path vts1  
  assumes p2 = make-polygonal-path vts2  
  assumes vts1 = take i vts  
  assumes vts2 = drop (i-1) vts  
  assumes n = length vts  
  assumes 1 ≤ i ∧ i < n  
  assumes x = (2^(i-1) - 1)/(2^(i-1))  
  shows path-image p1 = p'{0..x} ∧ path-image p2 = p'{x..1}  
 $\langle proof \rangle$ 
```

```
lemma drop-i-is-loop-free:  
  fixes vts :: (real^2) list  
  assumes m = length vts  
  assumes i ≤ m - 2  
  assumes vts' = drop i vts  
  assumes p = make-polygonal-path vts  
  assumes p' = make-polygonal-path vts'  
  assumes loop-free p  
  shows loop-free p'  
 $\langle proof \rangle$ 
```

```
lemma joinpaths-tl-transform:  
  assumes f = (λx::real. 2*x - 1)  
  assumes pathfinish g1 = pathstart g2  
  assumes p = g1 +++ g2  
  assumes x ≥ 1/2  
  shows p x = g2 (f x)  
 $\langle proof \rangle$ 
```

```
lemma joinpaths-tl-image-transform:  
  assumes f = (λx::real. 2*x - 1)  
  assumes pathfinish g1 = pathstart g2  
  assumes p = g1 +++ g2  
  assumes 1/2 ≤ a ∧ a ≤ b  
  shows p'{a..b} = g2'{f a..f b}  
 $\langle proof \rangle$ 
```

```
lemma vts-sublist-path-image:  
  assumes p = make-polygonal-path vts  
  assumes p' = make-polygonal-path vts'  
  assumes vts' = take j (drop i vts)  
  assumes m = length vts  
  assumes n = length vts'  
  assumes k = i + j  
  assumes k ≤ m - 1 ∧ 2 ≤ j
```

```

assumes  $x1 = (2^{\hat{i}} - 1)/(2^{\hat{i}})$ 
assumes  $x2 = (2^{(k-1)} - 1)/(2^{(k-1)})$ 
shows path-image  $p' = p^{\{x1..x2\}}$ 
⟨proof⟩

```

```

lemma one-append-simple-path:
fixes vts :: ( $\text{real}^2$ ) list
assumes vts = vts' @ [z]
assumes n = length vts
assumes n ≥ 3
assumes p = make-polygonal-path vts
assumes p' = make-polygonal-path vts'
assumes simple-path p
shows simple-path p'
⟨proof⟩

```

```

lemma take-i-is-loop-free:
fixes vts :: ( $\text{real}^2$ ) list
assumes n = length vts
assumes 2 ≤ i ∧ i ≤ n
assumes vts' = take i vts
assumes p = make-polygonal-path vts
assumes p' = make-polygonal-path vts'
assumes loop-free p
shows loop-free p'
⟨proof⟩

```

```

lemma sublist-is-loop-free:
fixes vts :: ( $\text{real}^2$ ) list
assumes p = make-polygonal-path vts
assumes p' = make-polygonal-path vts'
assumes loop-free p
assumes m = length vts
assumes n = length vts'
assumes sublist vts' vts
assumes n ≥ 2 ∧ m ≥ 2
shows loop-free p'
⟨proof⟩

```

```

lemma diff-points-path-image-set-property:
fixes a b::  $\text{real}^2$ 
assumes a ≠ b
shows path-image (linepath a b) ≠ {a, b}
⟨proof⟩

```

```

lemma polygonal-path-vertex-t:
assumes p = make-polygonal-path vts
assumes n = length vts
assumes n ≥ 1

```

```

assumes  $0 \leq i \wedge i < n - 1$ 
assumes  $x = (2^{\hat{i}} - 1)/(2^{\hat{i}})$ 
shows  $vts!i = p x$ 
⟨proof⟩

lemma loop-free-split-int:
assumes  $p = \text{make-polygonal-path } vts \wedge \text{loop-free } p$ 
assumes  $vts1 = \text{take } i \ vts$ 
assumes  $vts2 = \text{drop } (i-1) \ vts$ 
assumes  $c1 = \text{make-polygonal-path } vts1$ 
assumes  $c2 = \text{make-polygonal-path } vts2$ 
assumes  $n = \text{length } vts$ 
assumes  $1 \leq i \wedge i < n$ 
shows  $(\text{path-image } c1) \cap (\text{path-image } c2) \subseteq \{\text{pathstart } c1, \text{pathstart } c2\}$ 
(is  $?C1 \cap ?C2 \subseteq \{\text{pathstart } c1, \text{pathstart } c2\}$ )
⟨proof⟩

lemma loop-free-arc-split-int:
assumes  $p = \text{make-polygonal-path } vts \wedge \text{loop-free } p \wedge \text{arc } p$ 
assumes  $vts1 = \text{take } i \ vts$ 
assumes  $vts2 = \text{drop } (i-1) \ vts$ 
assumes  $c1 = \text{make-polygonal-path } vts1$ 
assumes  $c2 = \text{make-polygonal-path } vts2$ 
assumes  $n = \text{length } vts$ 
assumes  $1 \leq i \wedge i < n$ 
shows  $(\text{path-image } c1) \cap (\text{path-image } c2) \subseteq \{\text{pathstart } c2\}$ 
(is  $?C1 \cap ?C2 \subseteq \{\text{pathstart } c2\}$ )
⟨proof⟩

lemma loop-free-append:
assumes  $p = \text{make-polygonal-path } vts$ 
assumes  $p1 = \text{make-polygonal-path } vts1$ 
assumes  $p2 = \text{make-polygonal-path } vts2$ 
assumes  $vts = vts1 @ (tl vts2)$ 
assumes  $\text{loop-free } p1 \wedge \text{loop-free } p2$ 
assumes  $\text{path-image } p1 \cap \text{path-image } p2 \subseteq \{\text{pathstart } p1, \text{pathstart } p2\}$ 
assumes  $\text{last } vts2 \neq \text{hd } vts1 \longrightarrow \text{path-image } p1 \cap \text{path-image } p2 \subseteq \{\text{pathstart } p2\}$ 
assumes  $\text{last } vts1 = \text{hd } vts2$ 
assumes  $\text{arc } p1 \wedge \text{arc } p2$ 
shows  $\text{loop-free } p$ 
⟨proof⟩

lemma sublist-path-image-subset:
assumes  $\text{sublist } vts1 \ vts2$ 
assumes  $\text{length } vts1 \geq 1$ 
shows  $\text{path-image } (\text{make-polygonal-path } vts1) \subseteq \text{path-image } (\text{make-polygonal-path } vts2)$ 
⟨proof⟩

```

```

lemma integral-on-edge-subset-integral-on-path:
  assumes p = make-polygonal-path vts and
    (i:int) ∈ {0..<((length vts) − 1)} and
    x = vts!i and
    y = vts!(i+1)
  shows {v. integral-vec v ∧ v ∈ path-image (linepath x y)}
    ⊆ {v. integral-vec v ∧ v ∈ path-image p}
  ⟨proof⟩

lemma sublist-pair-integral-subset-integral-on-path:
  assumes p = make-polygonal-path vts and
    sublist [x, y] vts
  shows {v. integral-vec v ∧ v ∈ path-image (linepath x y)}
    ⊆ {v. integral-vec v ∧ v ∈ path-image p}
  ⟨proof⟩

lemma sublist-integral-subset-integral-on-path:
  assumes length ell ≥ 2
  assumes p = make-polygonal-path vts and
    sublist ell vts
  shows {v. integral-vec v ∧ v ∈ path-image (make-polygonal-path ell)}
    ⊆ {v. integral-vec v ∧ v ∈ path-image p}
  ⟨proof⟩

```

## 13 Reversing Polygonal Path Vertex List

```

lemma rev-vts-path-image:
  shows path-image (make-polygonal-path (rev vts)) = path-image (make-polygonal-path
vts)
  ⟨proof⟩

lemma rev-vts-is-loop-free:
  assumes p = make-polygonal-path vts
  assumes loop-free p
  shows loop-free (make-polygonal-path (rev vts))
  ⟨proof⟩

lemma rev-vts-is-polygon:
  assumes polygon-of p vts
  shows polygon (make-polygonal-path (rev vts))
  ⟨proof⟩

end
theory Linepath-Collinearity
imports Polygon-Lemmas

begin

```

## 14 Collinearity Properties

```

lemma points-on-linepath-collinear:
  assumes exists-c: ( $\exists c. a - b = c *_R u$ )
  assumes x-in-linepath:  $x \in \text{path-image}(\text{linepath } a \ b)$ 
  shows ( $\exists c. x - a = c *_R u$ ) ( $\exists c. b - x = c *_R u$ )
  ⟨proof⟩

lemma three-points-collinear-property:
  fixes a b:: real $^2$ 
  assumes exists-c1: ( $\exists c. a - x1 = c *_R u$ )
  assumes exists-c2: ( $\exists c. a - x2 = c *_R u$ )
  shows  $\exists c. x1 - x2 = c *_R u$ 
  ⟨proof⟩

lemma in-path-image-imp-collinear:
  fixes a b:: real $^2$ 
  assumes k ∈ path-image (linepath a b)
  shows collinear {a, b, k}
  ⟨proof⟩

lemma two-linepath-colinearity-property:
  fixes a b c d:: real $^2$ 
  assumes  $y \neq z \wedge \{y, z\} \subseteq (\text{path-image}(\text{linepath } a \ b)) \cap (\text{path-image}(\text{linepath } c \ d))$ 
  shows collinear {a, b, c, d}
  ⟨proof⟩

lemma polygon-vts-not-collinear:
  assumes polygon-of p vts
  shows  $\neg \text{collinear}(\text{set } vts)$ 
  ⟨proof⟩

lemma not-collinear-with-subset:
  assumes collinear A
  assumes  $\neg \text{collinear}(A \cup \{x\})$ 
  assumes card A > 2
  assumes a ∈ A
  shows  $\neg \text{collinear}((A - \{a\}) \cup \{x\})$ 
  ⟨proof⟩

lemma vec-diff-scale-collinear:
  fixes a b c :: real $^2$ 
  assumes  $b - a = m *_R (c - a)$ 
  shows collinear {a, b, c}
  ⟨proof⟩

```

## 15 Linepath Properties

**lemma** *good-linepath-comm*: *good-linepath a b vts*  $\implies$  *good-linepath b a vts*  
*(proof)*

**lemma** *finite-set-linepaths*:  
**assumes** *polygon*: *polygon p*  
**assumes** *polygonal-path*: *p = make-polygonal-path vts*  
**shows** *finite {(a, b). (a, b) ∈ set vts × set vts}*  
*(proof)*

**lemma** *linepaths-intersect-once-or-collinear*:  
**fixes** *a b c d :: real^2*  
**assumes** *path-image (linepath a b) ∩ path-image (linepath c d) ≠ {}*  
**shows** *collinear {a, b, c, d} ∨ (∃x. path-image (linepath a b) ∩ path-image (linepath c d) = {x})*  
*(proof)*

**lemma** *linepaths-intersect-once-or-collinear-alt*:  
**fixes** *a b c d :: real^2*  
**assumes** *path-image (linepath a b) ∩ path-image (linepath c d) ≠ {}*  
**shows** *collinear {a, b, c, d} ∨ card (path-image (linepath a b) ∩ path-image (linepath c d)) = 1*  
*(proof)*

**lemma** *path-image-linepath-union*:  
**fixes** *a b :: 'a::euclidean-space*  
**assumes** *d ∈ path-image (linepath a b)*  
**shows** *path-image (linepath a b) = path-image (linepath a d) ∪ path-image (linepath d b)*  
*(proof)*

**lemma** *path-image-linepath-split*:  
**assumes** *i < (length vts) - 1*  
**assumes** *x ∈ path-image (linepath (vts!i) (vts!(i+1)))*  
**assumes** *x-notin: x ∉ set vts*  
**shows** *path-image (make-polygonal-path vts) = path-image (make-polygonal-path ((take (i+1) vts) @ [x] @ (drop (i+1) vts)))*  
*(proof)*

**lemma** *linepath-split-is-loop-free*:  
**assumes** *d ∈ path-image (linepath a b)*  
**assumes** *d ∉ {a, b}*  
**shows** *loop-free (make-polygonal-path [a, d, b]) (is loop-free ?p)*  
*(proof)*

**lemma** *loop-free-linepath-split-is-loop-free*:  
**assumes** *p = make-polygonal-path vts*  
**assumes** *loop-free p*

```

assumes  $n = \text{length } vts$ 
assumes  $i < n - 1$ 
assumes  $x \in \text{path-image}(\text{linepath}(vts!i)(vts!(i+1))) \wedge x \notin \text{set } vts$ 
assumes  $vts' = (\text{take}(i+1)vts) @ [x] @ (\text{drop}(i+1)vts)$ 
assumes  $p' = \text{make-polygonal-path } vts'$ 
shows  $\text{loop-free } p' \wedge \text{path-image } p' = \text{path-image } p$ 
⟨proof⟩

```

```

lemma polygon-linepath-split-is-polygon:
assumes polygon-of  $p$   $vts$ 
assumes  $i < (\text{length } vts) - 1$ 
assumes  $a = vts!i \wedge b = vts!(i+1)$ 
assumes  $x \in \text{path-image}(\text{linepath } a \ b) \wedge x \notin \text{set } vts$ 
assumes  $vts' = (\text{take}(i+1)vts) @ [x] @ (\text{drop}(i+1)vts)$ 
shows polygon ( $\text{make-polygonal-path } vts'$ )
⟨proof⟩

```

## 16 Measure of linepaths

```

lemma linepath-is-negligible-vertical:
fixes  $a \ b :: \text{real}^2$ 
assumes  $a\$1 = b\$1$ 
defines  $p \equiv \text{linepath } a \ b$ 
shows negligible ( $\text{path-image } p$ )
⟨proof⟩

```

```

lemma linepath-is-negligible-non-vertical:
fixes  $a \ b :: \text{real}^2$ 
assumes  $a\$1 < b\$1$ 
defines  $p \equiv \text{linepath } a \ b$ 
shows negligible ( $\text{path-image } p$ )
⟨proof⟩

```

```

lemma linepath-is-negligible:
fixes  $a \ b :: \text{real}^2$ 
defines  $p \equiv \text{linepath } a \ b$ 
shows negligible ( $\text{path-image } p$ )
⟨proof⟩

```

```

lemma linepath-has-emeasure-0:
 $\text{emeasure lebesgue}(\text{path-image}(\text{linepath}(a::(\text{real}^2))(b::(\text{real}^2)))) = 0$ 
⟨proof⟩

```

```

lemma linepath-has-measure-0:
 $\text{measure lebesgue}(\text{path-image}(\text{linepath}(a::(\text{real}^2))(b::(\text{real}^2)))) = 0$ 
⟨proof⟩

```

end

```
theory Polygon-Convex-Lemmas
imports
```

*Polygon-Lemmas*  
*Linepath-Collinearity*

```
begin
```

## 17 Misc. Convex Polygon Properties

```
lemma polygon-path-image-subset-convex:
  assumes length vts > 0
  shows path-image (make-polygonal-path vts) ⊆ convex hull (set vts) (is path-image
?p ⊆ ?S)
  ⟨proof⟩

lemma convex-contains-simple-closed-path-imp-contains-path-inside:
  assumes convex S
  assumes simple-path p ∧ closed-path p
  assumes path-image p ⊆ S
  shows path-inside p ⊆ S
  ⟨proof⟩

lemma convex-polygon-is-convex-hull:
  assumes polygon p
  assumes convex (path-inside p ∪ path-image p)
  assumes p = make-polygonal-path vts
  shows convex hull (set vts) = path-inside p ∪ path-image p (is ?hull = ?poly)
  ⟨proof⟩

lemma convex-polygon-inside-is-convex-hull-interior:
  assumes polygon p
  assumes convex (path-inside p)
  assumes p = make-polygonal-path vts
  shows interior (convex hull (set vts)) = path-inside p
  ⟨proof⟩

lemma convex-polygon-inside-is-convex-hull-interior2:
  assumes polygon p
  assumes convex (path-inside p ∪ path-image p)
  assumes p = make-polygonal-path vts
  shows interior (convex hull (set vts)) = path-inside p
  ⟨proof⟩

lemma polygon-convex-iff:
  assumes polygon p
  shows convex (path-inside p) ←→ convex (path-inside p ∪ path-image p)
  ⟨proof⟩

lemma convex-polygon-frontier-is-path-image:
```

```

assumes polygon-of p vts
assumes convex (path-inside p)
shows frontier (convex hull (set vts)) = path-image p
⟨proof⟩

lemma convex-polygon-frontier-is-path-image2:
assumes polygon p
assumes convex (path-inside p)
shows frontier (path-image p ∪ path-inside p) = path-image p
⟨proof⟩

lemma convex-polygon-frontier-is-path-image3:
assumes polygon p
assumes convex (path-image p ∪ path-inside p)
shows frontier (path-image p ∪ path-inside p) = path-image p
⟨proof⟩

lemma polygon-frontier-is-path-image:
assumes polygon p
shows frontier (path-inside p) = path-image p
⟨proof⟩

lemma convex-path-inside-means-convex-polygon:
assumes polygon p
assumes frontier (convex hull (set vts)) = path-image p
shows convex (path-inside p)
⟨proof⟩

lemma convex-hull-of-polygon-is-convex-hull-of-vts:
assumes polygon-of p vts
shows convex hull (path-image p ∪ path-inside p) = convex hull (set vts)
⟨proof⟩

lemma convex-hull-frontier-polygon:
assumes polygon-of p vts
assumes ¬ set vts ⊆ frontier (convex hull (set vts))
shows ¬ convex (path-inside p)
⟨proof⟩

lemma frontier-int-subset:
assumes A ⊆ B
shows (frontier B) ∩ A ⊆ frontier A
⟨proof⟩

lemma in-frontier-in-subset:
assumes A ⊆ B
assumes x ∈ frontier B
assumes x ∈ A
shows x ∈ frontier A

```

$\langle proof \rangle$

**lemma** *in-frontier-in-subset-convex-hull*:

**assumes**  $A \subseteq B$

**assumes**  $x \in \text{frontier}(\text{convex hull } B)$

**assumes**  $x \in \text{convex hull } A$

**shows**  $x \in \text{frontier}(\text{convex hull } A)$

$\langle proof \rangle$

**lemma** *convex-hull-two-extreme-points*:

**fixes**  $S :: 'a::\text{euclidean-space set}$

**assumes**  $\text{finite } S$

**assumes**  $\text{convex hull } S \neq \{\}$

**assumes**  $\forall x. \text{convex hull } S \neq \{x\}$

**shows**  $\text{card}\{\{x. x \text{ extreme-point-of } (\text{convex hull } S)\}\} \geq 2$  (**is**  $\text{card } ?ep \geq 2$ )

$\langle proof \rangle$

**lemma** *convex-hull-two-vts-on-frontier*:

**fixes**  $S :: 'a::\text{euclidean-space set}$

**assumes**  $\text{card } S \geq 2$

**shows**  $\text{card}(S \cap \text{frontier}(\text{convex hull } S)) \geq 2$

$\langle proof \rangle$

## 18 Vertices on Convex Frontier Implies Polygon is Convex

**lemma** *convex-cut-aux*:

**assumes**  $\forall v \in S. z \cdot v \leq 0$

**shows**  $\text{convex hull } S \subseteq \{x. z \cdot x \leq 0\}$

$\langle proof \rangle$

**lemma** *convex-cut-aux'*:

**assumes**  $\forall v \in S. z \cdot v \geq 0$

**shows**  $\text{convex hull } S \subseteq \{x. z \cdot x \geq 0\}$

$\langle proof \rangle$

**lemma** *convex-cut*:

**assumes**  $z \neq 0$

**assumes**  $\{x. z \cdot x = 0\} \cap \text{interior}(\text{convex hull } S) \neq \{\}$

**obtains**  $v1 \ v2$  **where**  $v1 \neq v2 \wedge \{v1, v2\} \subseteq S \wedge v1 \in \{x. z \cdot x < 0\} \wedge v2 \in \{x. z \cdot x > 0\}$

$\langle proof \rangle$

**lemma** *affine-2-int-convex*:

**fixes**  $S :: 'a::\text{euclidean-space set}$

**assumes**  $\{a, b\} \subseteq S$

**assumes**  $\{a, b\} \subseteq \text{frontier}(\text{convex hull } S)$

**assumes**  $\text{affine hull } \{a, b\} \cap \text{interior}(\text{convex hull } S) \neq \{\}$

**shows**  $\text{affine hull } \{a, b\} \cap \text{convex hull } S = \text{convex hull } \{a, b\}$   
 $\langle proof \rangle$

**lemma** *halfplane-frontier-affine-hull*:

**fixes**  $b v :: \text{real}^2$   
**assumes**  $b \neq 0$   
**assumes**  $v \neq 0$   
**assumes**  $b \in \{x. v \cdot x = 0\}$   
**shows**  $\{x. v \cdot x = 0\} = \text{affine hull } \{0, b\}$   
 $\langle proof \rangle$

**lemma** *vts-on-convex-frontier-aux*:

**assumes**  $\text{polygon-of } p \text{ vts}$   
**assumes**  $\text{vts}!0 = 0$   
**assumes**  $\text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
**shows**  $\text{path-image } (\text{linepath } (\text{vts}!0) (\text{vts}!1)) \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
 $\langle proof \rangle$

**lemma** *vts-on-convex-frontier-aux'*:

**assumes**  $\text{polygon-of } p \text{ vts}$   
**assumes**  $\text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
**shows**  $\text{path-image } (\text{linepath } (\text{vts}!0) (\text{vts}!1)) \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
 $\langle proof \rangle$

**lemma** *vts-on-convex-frontier*:

**assumes**  $\text{polygon-of } p \text{ vts}$   
**assumes**  $\text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
**assumes**  $i < \text{length vts} - 1$   
**shows**  $\text{path-image } (\text{linepath } (\text{vts}!i) (\text{vts}!(i+1))) \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
 $\langle proof \rangle$

**lemma** *vts-on-frontier-means-path-image-on-frontier*:

**assumes**  $\text{polygon-of } p \text{ vts}$   
**assumes**  $\text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
**shows**  $\text{path-image } p \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
 $\langle proof \rangle$

**lemma** *vts-on-convex-frontier-interior*:

**assumes**  $\text{polygon-of } p \text{ vts}$   
**assumes**  $\text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
**shows**  $\text{path-inside } p = \text{interior } (\text{convex hull } (\text{set vts}))$   
 $\langle proof \rangle$

**lemma** *vts-subset-frontier*:

**assumes**  $\text{polygon-of } p \text{ vts}$   
**assumes**  $\text{set vts} \subseteq \text{frontier } (\text{convex hull } (\text{set vts}))$   
**shows**  $\text{convex } (\text{path-image } p \cup \text{path-inside } p)$   
 $\langle proof \rangle$

```

lemma convex-hull-of-nonconvex-polygon-strict-subset-ep:
  assumes polygon-of p vts
  assumes  $\neg (\text{convex}(\text{path-image } p \cup \text{path-inside } p))$ 
  shows  $\{v. v \text{ extreme-point-of } (\text{convex hull}(\text{set vts}))\} \subset \text{set vts}$ 
   $\langle \text{proof} \rangle$ 

lemma convex-hull-of-nonconvex-polygon-strict-subset:
  assumes polygon-of p vts
  assumes  $\neg (\text{convex}(\text{path-image } p \cup \text{path-inside } p))$ 
  shows  $\exists v \in \text{set vts}. v \in \text{interior}(\text{convex hull}(\text{set vts}))$ 
   $\langle \text{proof} \rangle$ 

lemma convex-polygon-means-linepaths-inside:
  fixes p :: R-to-R2
  assumes polygon-of p vts
  assumes convex-is:  $\text{convex hull}(\text{set vts}) = (\text{path-inside } p \cup \text{path-image } p)$ 
  assumes a-in:  $a \in (\text{path-inside } p \cup \text{path-image } p)$ 
  assumes b-in:  $b \in (\text{path-inside } p \cup \text{path-image } p)$ 
  shows  $\text{path-image}(\text{linepath } a b) \subseteq (\text{path-inside } p \cup \text{path-image } p)$ 
   $\langle \text{proof} \rangle$ 

end
theory Polygon-Splitting
imports
  HOL-Analysis.Complete-Measure
  Polygon-Jordan-Curve
  Polygon-Convex-Lemmas
begin

```

## 19 Polygon Splitting

```

lemma split-up-a-list-into-3-parts:
  fixes i j:: nat
  assumes  $i < \text{length vts} \wedge j < \text{length vts} \wedge i < j$ 
  shows
     $vts = (\text{take } i \text{ vts}) @ ((\text{vts} ! i) \# ((\text{take } (j - i - 1) (\text{drop } (\text{Suc } i) \text{ vts})) @ (\text{vts} ! j) \# \text{drop } (j - i) (\text{drop } (\text{Suc } i) \text{ vts})))$ 
   $\langle \text{proof} \rangle$ 

```

```

definition is-polygon-cut :: (real^2) list  $\Rightarrow$  real^2  $\Rightarrow$  real^2  $\Rightarrow$  bool where
  is-polygon-cut vts x y =
     $(x \neq y \wedge$ 
      $\text{polygon}(\text{make-polygonal-path } vts) \wedge$ 
      $\{x, y\} \subseteq \text{set vts} \wedge$ 
      $\text{path-image}(\text{linepath } x y) \cap \text{path-image}(\text{make-polygonal-path } vts) = \{x, y\} \wedge$ 
      $\text{path-image}(\text{linepath } x y) \cap \text{path-inside}(\text{make-polygonal-path } vts) \neq \{\})$ 

```

```

definition is-polygon-cut-path :: (real^2) list  $\Rightarrow$  R-to-R2  $\Rightarrow$  bool where

```

```

is-polygon-cut-path vts cutpath =
  (let x = pathstart cutpath ; y = pathfinish cutpath in
    (x ≠ y ∧
      polygon (make-polygonal-path vts) ∧
      {x, y} ⊆ set vts ∧
      simple-path cutpath ∧
      path-image cutpath ∩ path-image (make-polygonal-path vts) = {x, y} ∧
      path-image cutpath ∩ path-inside (make-polygonal-path vts) ≠ {}))

```

```

definition is-polygon-split :: (real^2) list ⇒ nat ⇒ nat ⇒ bool where
  is-polygon-split vts i j =
    (i < length vts ∧ j < length vts ∧ i < j ∧
     (let vts1 = (take i vts) in
      let vts2 = (take (j - i - 1) (drop (Suc i) vts)) in
      let vts3 = drop (j - i) (drop (Suc i) vts) in
      let x = vts ! i in
      let y = vts ! j in
      let p = make-polygonal-path (vts@[vts!0]) in
      let p1 = make-polygonal-path (x#(vts2@[y, x])) in
      let p2 = make-polygonal-path (vts1 @ [x, y] @ vts3 @[vts ! 0]) in
      let c1 = make-polygonal-path (x#(vts2@[y])) in
      let c2 = make-polygonal-path (vts1 @ [x, y] @ vts3) in
        (is-polygon-cut (vts@[vts!0]) x y ∧
         polygon p ∧ polygon p1 ∧ polygon p2 ∧
         path-inside p1 ∩ path-inside p2 = {} ∧
         path-inside p1 ∪ path-inside p2 ∪ (path-image (linepath x y) - {x, y}) =
          path-inside p
         ∧ ((path-image p1) - (path-image (linepath x y))) ∩ ((path-image p2) -
          (path-image (linepath x y)))
         = {}
         ∧ path-image p
         = ((path-image p1) - (path-image (linepath x y))) ∪ ((path-image p2) -
          (path-image (linepath x y))) ∪ {x, y}
        ))))

```

```

definition is-polygon-split-path :: (real^2) list ⇒ nat ⇒ nat ⇒ (real^2) list ⇒
  bool where
  is-polygon-split-path vts i j cutvts =
    (i < length vts ∧ j < length vts ∧ i < j ∧
     (let vts1 = (take i vts) in
      let vts2 = (take (j - i - 1) (drop (Suc i) vts)) in
      let vts3 = drop (j - i) (drop (Suc i) vts) in
      let x = vts!i in
      let y = vts!j in
      let cutpath = make-polygonal-path (x # cutvts @ [y]) in
      let p = make-polygonal-path (vts@[vts!0]) in
      let p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x])) in

```

```

let p2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @ [vts ! 0]) in
let c1 = make-polygonal-path (x#(vts2@[y])) in
let c2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3) in
(is-polygon-cut-path (vts@[vts!0])) cutpath ∧
polygon p ∧ polygon p1 ∧ polygon p2 ∧
path-inside p1 ∩ path-inside p2 = {} ∧
path-inside p1 ∪ path-inside p2 ∪ (path-image cutpath − {x, y}) = path-inside
p
∧ ((path-image p1) − (path-image cutpath)) ∩ ((path-image p2) − (path-image
cutpath)) = {}
∧ path-image p
= ((path-image p1) − (path-image cutpath)) ∪ ((path-image p2) − (path-image
cutpath)) ∪ {x, y}
)))

```

**lemma** polygon-split-add-measure:

**fixes**  $p\ p1\ p2 :: R\text{-to-}R2$

**assumes** is-polygon-split  $vts\ i\ j$

**assumes**  $vts1 = (\text{take } i \ vts)$

$vts2 = (\text{take } (j - i - 1) \ (\text{drop } (\text{Suc } i) \ vts))$

$vts3 = \text{drop } (j - i) \ (\text{drop } (\text{Suc } i) \ vts)$

$x = vts ! i$

$y = vts ! j$

$p = \text{make-polygonal-path } (vts@[vts!0])$

$p1 = \text{make-polygonal-path } (x#(vts2@[y, x]))$

$p2 = \text{make-polygonal-path } (vts1 @ [x, y] @ vts3 @ [vts ! 0])$

**defines**  $M1 \equiv \text{measure lebesgue } (\text{path-inside } p1)$  **and**

$M2 \equiv \text{measure lebesgue } (\text{path-inside } p2)$  **and**

$M \equiv \text{measure lebesgue } (\text{path-inside } p)$

**shows**  $M1 + M2 = M$

$\langle \text{proof} \rangle$

**lemma** polygonal-paths-measurable:

**shows** path-image (make-polygonal-path  $vts$ )  $\in$  sets lebesgue

$\langle \text{proof} \rangle$

**lemma** polygonal-path-has-emeasure-0:

**shows** emeasure lebesgue (path-image (make-polygonal-path  $vts$ )) = 0

$\langle \text{proof} \rangle$

**lemma** polygon-split-path-add-measure:

**fixes**  $p\ p1\ p2 :: R\text{-to-}R2$

**assumes** is-polygon-split-path  $vts\ i\ j\ cutvts$

**assumes**  $vts1 = (\text{take } i \ vts)$

$vts2 = (\text{take } (j - i - 1) \ (\text{drop } (\text{Suc } i) \ vts))$

$vts3 = \text{drop } (j - i) \ (\text{drop } (\text{Suc } i) \ vts)$

$x = vts ! i$

$y = vts ! j$

$p = \text{make-polygonal-path } (vts@[vts!0])$

```

 $p1 = \text{make-polygonal-path } (x\#(vts2 @ [y] @ (\text{rev } cutvts) @ [x]))$ 
 $p2 = \text{make-polygonal-path } (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @ [vts ! 0])$ 
defines  $M1 \equiv \text{measure lebesgue } (\text{path-inside } p1)$  and
     $M2 \equiv \text{measure lebesgue } (\text{path-inside } p2)$  and
     $M \equiv \text{measure lebesgue } (\text{path-inside } p)$ 
shows  $M1 + M2 = M$ 
⟨proof⟩

lemma polygon-cut-path-to-split-path-vtx0:
fixes  $p :: R\text{-to-}R2$ 
assumes polygon-p: polygon p and
     $i\text{-gt: } i > 0$  and
     $i\text{-lt: } i < \text{length } vts$  and
     $p\text{-is: } p = \text{make-polygonal-path } (vts @ [vts ! 0])$  and
     $\text{cutpath: } cutpath = \text{make-polygonal-path } ([vts!0] @ cutvts @ [vts!i])$  and
     $\text{have-cut: } \text{is-polygon-cut-path } (vts @ [vts!0]) \text{ cutpath}$ 
shows is-polygon-split-path vts 0 i cutvts
⟨proof⟩

lemma polygon-cut-path-to-split-path:
fixes  $p :: R\text{-to-}R2$ 
assumes polygon p
 $p = \text{make-polygonal-path } (vts @ [vts ! 0])$ 
 $\text{is-polygon-cut-path } (vts @ [vts!0]) \text{ cutpath}$ 
 $vts1 \equiv (\text{take } i \text{ vts})$ 
 $vts2 \equiv (\text{take } (j - i - 1) \text{ (drop } (\text{Suc } i) \text{ vts}))$ 
 $vts3 \equiv \text{drop } (j - i) \text{ (drop } (\text{Suc } i) \text{ vts)}$ 
 $x \equiv vts ! i$ 
 $y \equiv vts ! j$ 
 $\text{cutpath} = \text{make-polygonal-path } ([x] @ cutvts @ [y])$ 
 $i < \text{length } vts \wedge j < \text{length } vts \wedge i < j$ 
 $p1 \equiv \text{make-polygonal-path } (x\#(vts2 @ ([y] @ (\text{rev } cutvts) @ [x])))$  and
 $p2 \equiv \text{make-polygonal-path } (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @ [(vts1 @ [x]) ! 0])$ 
shows is-polygon-split-path vts i j cutvts
⟨proof⟩

lemma good-polygonal-path-implies-polygon-split-path:
assumes polygon p
assumes  $p = \text{make-polygonal-path } (vts @ [vts!0])$ 
assumes good-polygonal-path v1 cutvts v2  $(vts @ [vts!0])$ 
assumes  $i < \text{length } vts \wedge j < \text{length } vts$ 
assumes  $vts ! i = v1$ 
assumes  $vts ! j = v2$ 
assumes  $i < j$ 
shows is-polygon-split-path vts i j cutvts
⟨proof⟩

```

```

lemma good-path-iff:
  good-linepath a b vts  $\longleftrightarrow$  good-polygonal-path a [] b vts
   $\langle proof \rangle$ 

lemma polygon-cut-iff: is-polygon-cut (vts @ [vts!0]) (vts!i) (vts!j)
   $\longleftrightarrow$  is-polygon-cut-path (vts @ [vts!0]) (linepath (vts!i) (vts!j))
   $\langle proof \rangle$ 

lemma polygon-split-iff: is-polygon-split vts i j  $\longleftrightarrow$  is-polygon-split-path vts i j []
   $\langle proof \rangle$ 

lemma polygon-cut-to-split-vtx0:
  fixes p :: R-to-R2
  assumes polygon-p: polygon p and
    i-gt: i > 0 and
    i-lt: i < length vts and
    p-is: p = make-polygonal-path (vts @ [vts ! 0]) and
    have-cut: is-polygon-cut (vts @ [vts!0]) (vts!0) (vts!i)
  shows is-polygon-split vts 0 i
   $\langle proof \rangle$ 

lemma polygon-cut-to-split:
  fixes p :: R-to-R2
  assumes is-polygon-cut (vts @ [vts!0]) (vts!i) (vts!j)
    i < length vts  $\wedge$  j < length vts  $\wedge$  i < j
  shows is-polygon-split vts i j
   $\langle proof \rangle$ 

lemma good-linepath-implies-polygon-split:
  assumes polygon p
  assumes p = make-polygonal-path (vts @ [vts!0])
  assumes good-linepath v1 v2 (vts @ [vts!0])
  assumes i < length vts  $\wedge$  j < length vts
  assumes vts ! i = v1
  assumes vts ! j = v2
  assumes i < j
  shows is-polygon-split vts i j
   $\langle proof \rangle$ 

end
theory Triangle-Lemmas
imports
  Polygon-Convex-Lemmas
  Integral-Matrix
  Affine-Arithmetic.Floatairth-Expression
  HOL-Analysis.Topology-Euclidean-Space
  HOL-Analysis.Equivalence-Lebesgue-Henstock-Integration
  HOL-Analysis.Inner-Product

```

*HOL*-Analysis.Line-Segment  
*HOL*-Analysis.Convex-Euclidean-Space  
*HOL*-Analysis.Change-Of-Vars

**begin**

## 20 Triangles

**definition** elem-triangle ::  $\text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow \text{bool}$  **where**  
 $\text{elem-triangle } a\ b\ c \longleftrightarrow$   
 $\neg \text{collinear } \{a, b, c\}$   
 $\wedge \text{integral-vec } a \wedge \text{integral-vec } b \wedge \text{integral-vec } c$   
 $\wedge \{x. x \in \text{convex hull } \{a, b, c\} \wedge \text{integral-vec } x\} = \{a, b, c\}$

**definition** triangle-mat ::  $\text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2^{\wedge}2$  **where**  
 $\text{triangle-mat } a\ b\ c = \text{transpose } (\text{vector } [b - a, c - a])$

**definition** triangle-linear ::  $\text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow (\text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2)$  **where**  
 $\text{triangle-linear } a\ b\ c = (\lambda x. (\text{triangle-mat } a\ b\ c) *v x)$

**definition** triangle-affine ::  $\text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2 \Rightarrow (\text{real}^{\wedge}2 \Rightarrow \text{real}^{\wedge}2)$  **where**  
 $\text{triangle-affine } a\ b\ c = (\lambda x. a + (\text{triangle-mat } a\ b\ c) *v x)$

**abbreviation** unit-square  $\equiv$   
 $(\text{convex hull } \{\text{vector } [0, 0], \text{vector } [0, 1], \text{vector } [1, 1], \text{vector } [1, 0]\}) :: ((\text{real}^{\wedge}2) \text{ set})$

**abbreviation** unit-triangle  $\equiv$   
 $(\text{convex hull } \{\text{vector } [0, 0], \text{vector } [1, 0], \text{vector } [0, 1]\}) :: ((\text{real}^{\wedge}2) \text{ set})$

**abbreviation** unit-triangle'  $\equiv$   
 $(\text{convex hull } \{\text{vector } [1, 1], \text{vector } [1, 0], \text{vector } [0, 1]\}) :: ((\text{real}^{\wedge}2) \text{ set})$

**lemma** triangle-inside-is-convex-hull-interior:  
**assumes** polygon-of p [a, b, c, a]  
**shows** path-inside p = interior (convex hull {a, b, c})  
 $\langle \text{proof} \rangle$

**lemma** triangle-is-convex:  
**assumes** p = make-triangle a b c **and**  $\neg \text{collinear } \{a, b, c\}$   
**shows** convex (path-inside p) (**is convex**? s)  
 $\langle \text{proof} \rangle$

**lemma** affine-comp-linear-trans:  $\text{triangle-affine } a\ b\ c = (\lambda x. x + a) \circ (\text{triangle-linear } a\ b\ c)$   
 $\langle \text{proof} \rangle$

**lemma** triangle-linear-der:

```

fixes a b c :: real~2
defines T ≡ triangle-linear a b c
shows (T has-derivative T) (at x)
⟨proof⟩

lemma triangle-affine-der:
fixes a b c :: real~2
assumes S ∈ sets lebesgue and x ∈ S
defines A ≡ triangle-affine a b c and T ≡ triangle-linear a b c
shows x ∈ S ⇒ (A has-derivative T) (at x within S)
⟨proof⟩

lemma triangle-linear-inj:
fixes a b c :: real~2
assumes ¬ collinear {a, b, c}
defines L ≡ triangle-linear a b c
shows inj L
⟨proof⟩

lemma triangle-affine-inj:
fixes a b c :: real~2
assumes ¬ collinear {a, b, c}
defines A ≡ triangle-affine a b c
shows inj A
⟨proof⟩

lemma triangle-linear-integrable:
fixes a b c :: real~2
assumes S ∈ lmeasurable
defines T ≡ triangle-linear a b c
shows (λx. abs (det (matrix (T)))) integrable-on S (is (λx. ?c) integrable-on S)
⟨proof⟩

lemma measure-differentiable-image-eq-affine:
fixes a b c :: real~2
defines A ≡ triangle-affine a b c and T ≡ triangle-linear a b c
assumes S ∈ lmeasurable and ¬ collinear {a, b, c}
shows measure lebesgue (A ‘ S) = integral S (λx. abs (det (matrix T)))
⟨proof⟩

lemma triangle-affine-img:
fixes a b c :: real~2
defines A ≡ triangle-affine a b c
shows convex hull {a, b, c} = A ‘ unit-triangle
⟨proof⟩

lemma triangle-affine-e1-e2:
fixes a b c :: real~2
defines A ≡ triangle-affine a b c

```

```

shows (triangle-affine a b c) (vector [0, 0]) = a
    (triangle-affine a b c) (vector [1, 0]) = b
    (triangle-affine a b c) (vector [0, 1]) = c
⟨proof⟩

lemma triangle-measure-integral-of-det:
  fixes a b c :: real2
  defines S ≡ convex hull {a, b, c}
  assumes ¬ collinear {a, b, c}
  shows measure lebesgue S =
    integral unit-triangle (λ(x::real2). abs (det (matrix (triangle-linear a b
c))))
⟨proof⟩

lemma triangle-affine-preserves-interior:
  assumes A = triangle-affine a b c and L = triangle-linear a b c
  assumes ¬ collinear {a, b, c}
  shows A ` (interior S) = interior (A ` S)
⟨proof⟩

lemma triangle-affine-preserves-affine-hull:
  assumes A = triangle-affine a b c
  assumes ¬ collinear {a, b, c}
  shows A ` (affine hull S) = affine hull (A ` S)
⟨proof⟩

lemma triangle-measure-convex-hull-measure-path-inside-same:
  assumes p-triangle: p = make-triangle a b c
  assumes elem-triangle: elem-triangle a b c
  shows measure lebesgue (convex hull {a, b, c}) = measure lebesgue (path-inside
p)
  (is measure lebesgue ?S = measure lebesgue ?I)
⟨proof⟩

lemma on-triangle-path-image-cases:
  assumes p = make-triangle a b c
  assumes d ∈ path-image p
  shows d ∈ path-image (linopath a b) ∨ d ∈ path-image (linopath b c) ∨ d ∈
path-image (linopath c a)
⟨proof⟩

lemma on-triangle-frontier-cases:
  fixes a b c :: real2
  assumes ¬ collinear {a, b, c}
  assumes d ∈ frontier (convex hull {a, b, c})
  shows d ∈ path-image (linopath a b) ∨ d ∈ path-image (linopath b c) ∨ d ∈
path-image (linopath c a)
⟨proof⟩

```

```

lemma triangle-path-image-subset-convex:
  assumes p = make-triangle a b c
  shows path-image p ⊆ convex hull {a, b, c}
  ⟨proof⟩

lemma triangle-convex-hull:
  assumes p = make-triangle a b c and ¬ collinear {a, b, c}
  shows convex hull {a, b, c} = (path-image p) ∪ (path-inside p)
  ⟨proof⟩

end
theory Unit-Geometry
imports
  HOL-Analysis.Polytope
  Polygon-Jordan-Curve
  Triangle-Lemmas

begin

```

## 21 Measure Setup

```

lemma finite-convex-is-measurable:
  fixes p :: (real^2) set
  assumes p = convex hull l and finite l
  shows p ∈ sets lebesgue
  ⟨proof⟩

lemma unit-square-lebesgue: unit-square ∈ sets lebesgue
  ⟨proof⟩

lemma unit-triangle-lebesgue: unit-triangle ∈ sets lebesgue
  ⟨proof⟩

lemma unit-triangle-lmeasurable: unit-triangle ∈ lmeasurable
  ⟨proof⟩

```

## 22 Unit Triangle

```

lemma unit-triangle-vts-not-collinear:
  ¬ collinear {(vector [0, 0])::real^2, vector [1, 0], vector [0, 1]}
  (is ¬ collinear {?a, ?b, ?c})
  ⟨proof⟩

lemma unit-triangle-convex:
  assumes p = (make-polygonal-path [vector [0, 0], vector [1, 0], vector [0, 1],
  vector [0, 0]])
  (is p = make-polygonal-path [?O, ?e1, ?e2, ?O])

```

```

shows convex (path-inside p)
⟨proof⟩

lemma unit-triangle-char:
shows unit-triangle = {x. 0 ≤ x $ 1 ∧ 0 ≤ x $ 2 ∧ x $ 1 + x $ 2 ≤ 1}
  (is unit-triangle = ?S)
⟨proof⟩

lemma unit-triangle-interior-char:
shows interior unit-triangle = {x. 0 < x $ 1 ∧ 0 < x $ 2 ∧ x $ 1 + x $ 2 <
  1}
  (is interior unit-triangle = ?S)
⟨proof⟩

lemma unit-triangle-is-elementary: elem-triangle (vector [0, 0]) (vector [1, 0])
(vectors [0, 1])
  (is elem-triangle ?a ?b ?c)
⟨proof⟩

lemma unit-triangles-same-area:
measure lebesgue unit-triangle' = measure lebesgue unit-triangle
⟨proof⟩

```

## 23 Unit Square

```

lemma convex-hull-4:
shows convex hull {a,b,c,d} = { u *R a + v *R b + w *R c + t *R d | u v w t. 0 ≤ u
  ∧ 0 ≤ v ∧ 0 ≤ w ∧ 0 ≤ t ∧ u + v + w + t = 1}
⟨proof⟩

lemma unit-square-characterization-helper:
fixes a b :: real
assumes 0 ≤ a ∧ a ≤ 1 ∧ 0 ≤ b ∧ b ≤ 1 and
  a ≤ b
obtains u v w t where
  vector [a, b] = u *R ((vector [0, 0])::real2)
    + v *R (vector [0, 1])
    + w *R (vector [1, 1])
    + t *R (vector [1, 0])
  ∧ 0 ≤ u ∧ 0 ≤ v ∧ 0 ≤ w ∧ 0 ≤ t ∧ u + v + w + t = 1
⟨proof⟩

lemma unit-square-characterization:
shows unit-square = {x. 0 ≤ x$1 ∧ x$1 ≤ 1 ∧ 0 ≤ x$2 ∧ x$2 ≤ 1} (is unit-square
= ?S)
⟨proof⟩

```

```

lemma e1e2-basis:
defines e1 ≡ (vector [1, 0])::(real2) and

```

```

 $e2 \equiv (\text{vector } [0, 1])::(\text{real}^2)$ 
shows  $e1 = \text{axis } 1 \ (1::\text{real})$  and  $e1 \in (\text{Basis}::((\text{real}^2) \ \text{set}))$  and
 $e2 = \text{axis } 2 \ (1::\text{real})$  and  $e2 \in (\text{Basis}::((\text{real}^2) \ \text{set}))$ 
⟨proof⟩

```

```

lemma unit-square-cbox:  $\text{unit-square} = \text{cbox } (\text{vector } [0, 0]) \ (\text{vector } [1, 1])$ 
⟨proof⟩

```

```

lemma unit-square-area:  $\text{measure lebesgue } \text{unit-square} = 1$ 
⟨proof⟩

```

## 24 Unit Triangle Area is 1/2

```

lemma unit-triangle'-char:
shows  $\text{unit-triangle}' = \{x. \ x \$ 1 \leq 1 \wedge x \$ 2 \leq 1 \wedge x \$ 1 + x \$ 2 \geq 1\}$ 
⟨proof⟩

```

```

lemma unit-square-split-diag:
shows  $\text{unit-square} = \text{unit-triangle} \cup \text{unit-triangle}'$ 
⟨proof⟩

```

```

lemma unit-triangle-INT-unit-triangle'-measure:
measure lebesgue ( $\text{unit-triangle} \cap \text{unit-triangle}'$ ) = 0
⟨proof⟩

```

```

lemma unit-triangle-area:  $\text{measure lebesgue } \text{unit-triangle} = 1/2$ 
⟨proof⟩

```

```

end
theory Elementary-Triangle-Area
imports
    Unit-Geometry

```

```

begin

```

## 25 Area of Elementary Triangle is 1/2

```

lemma nonint-in-square-img-IMP-nonint-triangle-img:
assumes  $A = \text{triangle-affine } a \ b \ c$ 
assumes  $x \in \text{unit-square}$ 
assumes  $\neg \text{integral-vec } x$ 
assumes  $\text{integral-vec } (A \ x)$ 
assumes  $\text{elem-triangle } a \ b \ c$ 
obtains  $x'$  where  $x' \in \text{unit-triangle} \wedge \neg \text{integral-vec } x' \wedge \text{integral-vec } (A \ x')$ 
⟨proof⟩

```

```

lemma elem-triangle-integral-mat-bij:
fixes  $a \ b \ c :: \text{real}^2$ 

```

```

assumes elem-triangle a b c
defines L ≡ triangle-mat a b c
shows integral-mat-bij L
⟨proof⟩

lemma elem-triangle-measure-integral-of-1:
  fixes a b c :: real^2
  defines S ≡ convex hull {a, b, c}
  assumes elem-triangle a b c
  shows measure lebesgue S = integral unit-triangle (λ(x::real^2). 1)
⟨proof⟩

lemma elem-triangle-area-is-half:
  fixes a b c :: real^2
  assumes elem-triangle a b c
  defines S ≡ convex hull {a, b, c}
  shows measure lebesgue S = 1/2 (is ?S-area = 1/2)
⟨proof⟩

end
theory Pick
imports
  Polygon-Splitting
  Elementary-Triangle-Area
begin

```

## 26 Setup

### 26.1 Integral Points Cardinality Properties

```

lemma bounded-finite:
  fixes A:: (real^2) set
  assumes bounded A
  shows finite {x::(real^2). integral-vec x ∧ x ∈ A} (is finite ?A-int)
⟨proof⟩

lemma finite-path-image:
  assumes polygon p
  shows finite {x. integral-vec x ∧ x ∈ path-image p}
⟨proof⟩

lemma finite-path-inside:
  assumes polygon p
  shows finite {x. integral-vec x ∧ x ∈ path-inside p}
⟨proof⟩

lemma bounded-finite-inside:
  fixes B:: (real^2) set
  assumes simple-path p

```

```

shows bounded (path-inside p)
⟨proof⟩

lemma finite-integral-points-path-image:
assumes simple-path p
shows finite {x. integral-vec x ∧ x ∈ path-image p}
⟨proof⟩

```

```

lemma finite-integral-points-path-inside:
assumes simple-path p
shows finite {x. integral-vec x ∧ x ∈ path-inside p}
⟨proof⟩

```

## 27 Pick splitting

```

lemma pick-split-path-union-main:
assumes is-split: is-polygon-split-path vts i j cutvts
assumes vts1 = (take i vts)
assumes vts2 = (take (j - i - 1) (drop (Suc i) vts))
assumes vts3 = drop (j - i) (drop (Suc i) vts)
assumes x = vts!i
assumes y = vts!j
assumes cutpath = make-polygonal-path (x # cutvts @ [y])
assumes p: p = make-polygonal-path (vts@[vts!0]) (is p = make-polygonal-path
?p-vts)
assumes p1: p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x]))
(is p1 = make-polygonal-path ?p1-vts)
assumes p2: p2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @
[vts ! 0]) (is p2 = make-polygonal-path ?p2-vts)
assumes I1: I1 = card {x. integral-vec x ∧ x ∈ path-inside p1}
assumes B1: B1 = card {x. integral-vec x ∧ x ∈ path-image p1}
assumes I2: I2 = card {x. integral-vec x ∧ x ∈ path-inside p2}
assumes B2: B2 = card {x. integral-vec x ∧ x ∈ path-image p2}
assumes I: I = card {x. integral-vec x ∧ x ∈ path-inside p}
assumes B: B = card {x. integral-vec x ∧ x ∈ path-image p}
assumes all-integral-vts: all-integral vts
shows measure lebesgue (path-inside p1) = I1 + B1/2 - 1
    ⇒ measure lebesgue (path-inside p2) = I2 + B2/2 - 1
    ⇒ measure lebesgue (path-inside p) = I + B/2 - 1
measure lebesgue (path-inside p) = I + B/2 - 1
    ⇒ measure lebesgue (path-inside p2) = I2 + B2/2 - 1
    ⇒ measure lebesgue (path-inside p1) = I1 + B1/2 - 1
measure lebesgue (path-inside p) = I + B/2 - 1
    ⇒ measure lebesgue (path-inside p1) = I1 + B1/2 - 1
    ⇒ measure lebesgue (path-inside p2) = I2 + B2/2 - 1
⟨proof⟩

```

```

lemma pick-split-union:
assumes is-split: is-polygon-split vts i j

```

```

assumes vts1 = (take i vts)
assumes vts2 = (take (j - i - 1) (drop (Suc i) vts))
assumes vts3 = drop (j - i) (drop (Suc i) vts)
assumes x = vts ! i
assumes y = vts ! j
assumes p: p = make-polygonal-path (vts@[vts!0]) (is p = make-polygonal-path
?p-vts)
assumes p1: p1 = make-polygonal-path (x#(vts2@[y, x])) (is p1 = make-polygonal-path
?p1-vts)
assumes p2: p2 = make-polygonal-path (vts1 @ [x, y] @ vts3 @ [vts ! 0]) (is p2
= make-polygonal-path ?p2-vts)
assumes I1: I1 = card {x. integral-vec x ∧ x ∈ path-inside p1}
assumes B1: B1 = card {x. integral-vec x ∧ x ∈ path-image p1}
assumes pick1: measure lebesgue (path-inside p1) = I1 + B1/2 - 1
assumes I2: I2 = card {x. integral-vec x ∧ x ∈ path-inside p2}
assumes B2: B2 = card {x. integral-vec x ∧ x ∈ path-image p2}
assumes pick2: measure lebesgue (path-inside p2) = I2 + B2/2 - 1
assumes I: I = card {x. integral-vec x ∧ x ∈ path-inside p}
assumes B: B = card {x. integral-vec x ∧ x ∈ path-image p}
assumes all-integral-vts: all-integral vts
shows measure lebesgue (path-inside p) = I + B/2 - 1
    measure lebesgue (path-inside p) = measure lebesgue (path-inside p1) +
    measure lebesgue (path-inside p2)
⟨proof⟩

```

```

lemma pick-split-path-union:
assumes is-split: is-polygon-split-path vts i j cutvts
assumes vts1 = (take i vts)
assumes vts2 = (take (j - i - 1) (drop (Suc i) vts))
assumes vts3 = drop (j - i) (drop (Suc i) vts)
assumes x = vts!i
assumes y = vts!j
assumes cutpath = make-polygonal-path (x # cutvts @ [y])
assumes p: p = make-polygonal-path (vts@[vts!0]) (is p = make-polygonal-path
?p-vts)
assumes p1: p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x]))
(is p1 = make-polygonal-path ?p1-vts)
assumes p2: p2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @
[vts ! 0]) (is p2 = make-polygonal-path ?p2-vts)
assumes I1: I1 = card {x. integral-vec x ∧ x ∈ path-inside p1}
assumes B1: B1 = card {x. integral-vec x ∧ x ∈ path-image p1}
assumes pick1: measure lebesgue (path-inside p1) = I1 + B1/2 - 1
assumes I2: I2 = card {x. integral-vec x ∧ x ∈ path-inside p2}
assumes B2: B2 = card {x. integral-vec x ∧ x ∈ path-image p2}
assumes pick2: measure lebesgue (path-inside p2) = I2 + B2/2 - 1
assumes I: I = card {x. integral-vec x ∧ x ∈ path-inside p}
assumes B: B = card {x. integral-vec x ∧ x ∈ path-image p}
assumes all-integral-vts: all-integral vts
shows measure lebesgue (path-inside p) = I + B/2 - 1

```

$\langle proof \rangle$

```

lemma pick-triangle-basic-split:
  assumes  $p = \text{make-triangle } a \ b \ c$  and  $\text{distinct } [a, b, c]$  and  $\neg \text{collinear } \{a, b, c\}$  and
     $d\text{-prop: } d \in \text{path-image}(\text{linepath } a \ b) \wedge d \notin \{a, b, c\}$ 
  shows  $\text{good-linepath } c \ d \ [a, d, b, c, a]$ 
     $\wedge \text{path-image}(\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p$ 
 $\langle proof \rangle$ 

```

## 28 Convex Hull Has Good Linepath

```

lemma leq-2-extreme-points-means-collinear:
  fixes vts :: 'a::euclidean-space set
  assumes finite vts
  assumes  $\text{card } \{v. v \text{ extreme-point-of } (\text{convex hull } vts)\} \leq 2$ 
  shows collinear vts
 $\langle proof \rangle$ 

```

```

lemma convex-hull-non-extreme-point-in-open-seg:
  assumes  $H = \text{convex hull } vts$ 
  assumes  $x \in H - \{v. v \text{ extreme-point-of } H\}$ 
  shows  $\exists a \ b. a \in H \wedge b \in H \wedge x \in \text{open-segment } a \ b$ 
 $\langle proof \rangle$ 

```

```

lemma convex-hull-extreme-points-vertex-split:
  fixes vts :: (real^2) set
  assumes  $H = \text{convex hull } vts$ 
  assumes finite vts
  assumes  $\text{card } \{v. v \text{ extreme-point-of } H\} \geq 4$ 
  assumes  $\{a, b, c\} \subseteq \{v. v \text{ extreme-point-of } H\} \wedge \text{distinct } [a, b, c]$ 
  shows  $\text{path-image}(\text{linepath } a \ b) \cap \text{interior } H \neq \{\}$ 
     $\vee \text{path-image}(\text{linepath } b \ c) \cap \text{interior } H \neq \{\}$ 
     $\vee \text{path-image}(\text{linepath } c \ a) \cap \text{interior } H \neq \{\}$ 
 $\langle proof \rangle$ 

```

```

lemma convex-hull-has-vertex-split-helper-wlog:
  assumes  $p = \text{make-triangle } a \ b \ c$  and  $\text{distinct } [a, b, c]$  and  $\neg \text{collinear } \{a, b, c\}$  and
     $d\text{-prop: } d \in \text{path-image}(\text{linepath } a \ b) \wedge d \notin \{a, b, c\}$ 
  shows  $\text{path-image}(\text{linepath } c \ d) \cap \text{path-inside } p \neq \{\}$ 
 $\langle proof \rangle$ 

```

```

lemma convex-hull-has-vertex-split-helper:
  assumes  $p = \text{make-triangle } a \ b \ c$  and  $\text{distinct } [a, b, c]$  and  $\neg \text{collinear } \{a, b, c\}$  and
     $d\text{-prop: } d \in \text{path-image } p \wedge d \notin \{a, b, c\}$ 
  shows  $\exists x \ y. \{x, y\} \subseteq \{a, b, c, d\} \wedge x \neq y \wedge \text{path-image}(\text{linepath } x \ y) \cap$ 
     $\text{path-inside } p \neq \{\}$ 

```

$\langle proof \rangle$

```
lemma convex-hull-has-vertex-split:
  fixes vts :: (real^2) set
  assumes H = convex hull vts
  assumes ¬ collinear vts
  assumes card vts > 3
  assumes finite vts
  shows ∃ a b. {a, b} ⊆ vts ∧ a ≠ b ∧ path-image (linepath a b) ∩ interior H ≠ {}
⟨proof⟩
```

```
lemma convex-polygon-has-good-linepath-helper:
  assumes polygon-of p vts
  assumes convex (path-inside p ∪ path-image p)
  assumes card (set vts) > 3
  obtains a b where {a, b} ⊆ set vts ∧ a ≠ b ∧ ¬ path-image (linepath a b) ⊆
    path-image p
⟨proof⟩
```

```
lemma convex-polygon-has-good-linepath:
  assumes convex (path-inside p ∪ path-image p)
  assumes polygon p
  assumes p = make-polygonal-path vts
  assumes card (set vts) > 3
  shows ∃ a b. good-linepath a b vts
⟨proof⟩
```

## 29 Pick's Theorem

```
definition integral-inside:
  integral-inside p = {x. integral-vec x ∧ x ∈ path-inside p}
```

```
definition integral-boundary:
  integral-boundary p = {x. integral-vec x ∧ x ∈ path-image p}
```

### 29.1 Pick's Theorem Triangle Case

```
definition pick-triangle:
  pick-triangle p a b c ↔
    p = make-triangle a b c
    ∧ all-integral [a, b, c]
    ∧ distinct [a, b, c]
    ∧ ¬ collinear {a, b, c}
```

```
definition pick-holds:
  pick-holds p ↔
    (let I = card {x. integral-vec x ∧ x ∈ path-inside p} in
     let B = card {x. integral-vec x ∧ x ∈ path-image p} in
```

*measure lebesgue (path-inside p) = I + B/2 - 1*

**lemma** *pick-triangle-wlog-helper*:  
**assumes** *pick-triangle p a b c and*  
*I = card (integral-inside p) and*  
*B = card (integral-boundary p) and*  
*integral-inside p = {} and*  
*integral-vec d ∧ d ∈ path-image (linepath a b) ∧ d ∉ {a, b, c} and d ∉ {a, b, c} and*  
*ih: ∏p' a' b' c'. (card (integral-inside p') + card (integral-boundary p') < I + B) ⇒ pick-triangle p' a' b' c' ⇒ pick-holds p'*  
**shows** *measure lebesgue (path-inside p) = I + B/2 - 1*  
*{proof}*

**lemma** *pick-triangle-helper*:  
**assumes** *pick-triangle p a b c and*  
*I = card (integral-inside p) and*  
*B = card (integral-boundary p) and*  
*integral-inside p = {} and*  
*integral-vec d ∧ d ∉ {a, b, c} and d ∉ {a, b, c} and*  
*d ∈ path-image (linepath a b)*  
*∨ d ∈ path-image (linepath b c)*  
*∨ d ∈ path-image (linepath c a) and*  
*ih: ∏p' a' b' c'. (card (integral-inside p') + card (integral-boundary p') < I + B) ⇒ pick-triangle p' a' b' c' ⇒ pick-holds p'*  
**shows** *measure lebesgue (path-inside p) = I + B/2 - 1*  
*{proof}*

**lemma** *triangle-3-split-helper*:  
**fixes** *a b :: 'a::euclidean-space*  
**assumes** *a ∈ frontier S*  
**assumes** *b ∈ interior S*  
**assumes** *convex S*  
**assumes** *closed S*  
**shows** *path-image (linepath a b) ∩ frontier S = {a}*  
*{proof}*

**lemma** *unit-triangle-interior-point-not-collinear-e1-e2*:  
**assumes** *p = make-triangle (vector [0, 0]) (vector [1, 0]) (vector [0, 1])*  
*(is p = make-triangle ?O ?e1 ?e2)*  
**assumes** *z ∈ path-inside p*  
**shows** *¬ collinear {?O, ?e1, z}*  
*{proof}*

**lemma** *triangle-interior-point-not-collinear-vertices-wlog-helper*:  
**assumes** *p = make-triangle a b c*  
**assumes** *polygon p*  
**assumes** *z ∈ path-inside p*  
**shows** *¬ collinear {a, b, z}*

$\langle proof \rangle$

**lemma** triangle-interior-point-not-collinear-vertices:  
  **assumes**  $p = \text{make-triangle } a \ b \ c$   
  **assumes** polygon  $p$   
  **assumes**  $z \in \text{path-inside } p$   
  **shows**  $\neg \text{collinear } \{a, b, z\} \wedge \neg \text{collinear } \{a, c, z\} \wedge \neg \text{collinear } \{b, c, z\}$   
 $\langle proof \rangle$

**lemma** triangle-3-split:  
  **assumes**  $p = \text{make-triangle } a \ b \ c$   
  **assumes** polygon  $p$   
  **assumes**  $z \in \text{path-inside } p$   
  **shows** is-polygon-split-path  $[a, b, c] \ 0 \ 1 \ [z]$   
    is-polygon-split  $[a, z, b, c] \ 1 \ 3$   
     $a \notin \text{path-image } (\text{make-triangle } z \ b \ c) \cup \text{path-inside } (\text{make-triangle } z \ b \ c)$   
     $b \notin \text{path-image } (\text{make-triangle } a \ z \ c) \cup \text{path-inside } (\text{make-triangle } a \ z \ c)$   
     $c \notin \text{path-image } (\text{make-triangle } a \ b \ z) \cup \text{path-inside } (\text{make-triangle } a \ b \ z)$   
 $\langle proof \rangle$

**lemma** smaller-triangle:  
  **assumes**  $\neg \text{collinear } \{a, b, c\} \wedge \neg \text{collinear } \{a', b', c'\}$   
  **assumes**  $p = \text{make-triangle } a \ b \ c$   
  **assumes**  $p' = \text{make-triangle } a' \ b' \ c'$   
  **assumes** path-inside  $p \subseteq \text{path-inside } p'$   
  **assumes**  $\exists d. \text{integral-vec } d \wedge d \in \text{path-image } p' \cup \text{path-inside } p' \wedge d \notin \text{path-image } p \cup \text{path-inside } p$   
  **shows** card (integral-inside  $p$ ) + card (integral-boundary  $p$ ) < card (integral-inside  $p'$ ) + card (integral-boundary  $p'$ )  
 $\langle proof \rangle$

**lemma** pick-elem-triangle:  
  **fixes**  $p :: R\text{-to-}R^2$   
  **assumes**  $p\text{-triangle}: p = \text{make-triangle } a \ b \ c$   
  **assumes** elem-triangle: elem-triangle  $a \ b \ c$   
  **assumes**  $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$  **and**  
     $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$   
  **shows** measure lebesgue (path-inside  $p$ ) =  $I + B/2 - 1$   
 $\langle proof \rangle$

**lemma** pick-triangle-lemma:  
  **fixes**  $p :: R\text{-to-}R^2$   
  **assumes**  $p = \text{make-triangle } a \ b \ c$  **and** all-integral  $[a, b, c]$  **and** distinct  $[a, b, c]$   
  **and**  $\neg \text{collinear } \{a, b, c\}$   
     $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$  **and**  
     $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$   
  **shows** measure lebesgue (path-inside  $p$ ) =  $I + B/2 - 1$   
 $\langle proof \rangle$

## 29.2 Pocket properties

```

definition index-not-in-set :: ( $\text{real}^2$ ) list  $\Rightarrow$  ( $\text{real}^2$ ) set  $\Rightarrow$  nat  $\Rightarrow$  bool
  where index-not-in-set vts A i  $\longleftrightarrow$  i  $\in \{i. i < \text{length } vts \wedge vts ! i \notin A\}$ 

definition min-index-not-in-set:: ( $\text{real}^2$ ) list  $\Rightarrow$  ( $\text{real}^2$ ) set  $\Rightarrow$  nat
  where min-index-not-in-set vts A = (LEAST i. index-not-in-set vts A i)

definition nonzero-index-in-set :: ( $\text{real}^2$ ) list  $\Rightarrow$  ( $\text{real}^2$ ) set  $\Rightarrow$  nat  $\Rightarrow$  bool
where
  nonzero-index-in-set vts A i  $\longleftrightarrow$  i  $\in \{i. 0 < i \wedge i < \text{length } vts \wedge vts ! i \in A\}$ 

definition min-nonzero-index-in-set :: ( $\text{real}^2$ ) list  $\Rightarrow$  ( $\text{real}^2$ ) set  $\Rightarrow$  nat where
  min-nonzero-index-in-set vts A = (LEAST i. nonzero-index-in-set vts A i)

definition construct-pocket-0 :: ( $\text{real}^2$ ) list  $\Rightarrow$  ( $\text{real}^2$ ) set  $\Rightarrow$  ( $\text{real}^2$ ) list where
  construct-pocket-0 vts A = take ((min-nonzero-index-in-set vts A) + 1) vts

definition is-pocket-0 :: ( $\text{real}^2$ ) list  $\Rightarrow$  ( $\text{real}^2$ ) list  $\Rightarrow$  bool where
  is-pocket-0 vts vts'  $\longleftrightarrow$ 
    polygon (make-polygonal-path vts)
     $\wedge (\exists i. vts' = \text{take } i vts)$ 
     $\wedge 3 \leq \text{length } vts' \wedge \text{length } vts' < \text{length } vts$ 
     $\wedge \text{hd } vts' \in \text{frontier } (\text{convex hull } (\text{set } vts)) \wedge \text{last } vts' \in \text{frontier } (\text{convex hull } (\text{set } vts))$ 
     $\wedge \text{set } (\text{tl } (\text{butlast } vts')) \subseteq \text{interior } (\text{convex hull } (\text{set } vts))$ 

definition fill-pocket-0 :: ( $\text{real}^2$ ) list  $\Rightarrow$  nat  $\Rightarrow$  ( $\text{real}^2$ ) list where
  fill-pocket-0 vts i = (hd vts) # (drop (i-1) vts)

lemma min-nonzero-index-in-set-exists:
  assumes set (tl vts)  $\cap A \neq \{\}$ 
  shows  $\exists i. \text{nonzero-index-in-set } vts A i$ 
  ⟨proof⟩

lemma min-nonzero-index-in-set-defined:
  assumes set (tl vts)  $\cap A \neq \{\}$ 
  defines i  $\equiv$  min-nonzero-index-in-set vts A
  shows nonzero-index-in-set vts A i  $\wedge (\forall j < i. \neg \text{nonzero-index-in-set } vts A j)$ 
  ⟨proof⟩

lemma min-index-not-in-set-exists:
  assumes set vts  $\supset A$ 
  shows  $\exists i. \text{index-not-in-set } vts A i$ 
  ⟨proof⟩

lemma min-index-not-in-set-defined:
  assumes set vts  $\supset A$ 
```

```

defines  $i \equiv \min\text{-index-not-in-set} vts A$ 
shows  $\text{index-not-in-set} vts A i \wedge (\forall j < i. \neg \text{index-not-in-set} vts A j)$ 
⟨proof⟩

lemma min-nonzero-index-in-set-bound:
assumes  $\text{set}(\text{tl } vts) \cap A \neq \{\}$ 
shows  $\min\text{-nonzero-index-in-set} vts A < \text{length } vts$ 
⟨proof⟩

lemma construct-pocket-0-subset-vts:
assumes  $\text{set}(\text{tl } vts) \cap A \neq \{\}$ 
shows  $\text{set}(\text{construct-pocket-0 } vts A) \subseteq \text{set } vts$ 
⟨proof⟩

lemma min-index-not-in-set-0:
assumes  $\text{set } vts \supset A$ 
assumes  $vts!0 \in A$ 
defines  $i \equiv \min\text{-index-not-in-set} vts A$ 
defines  $r \equiv i - 1$ 
shows  $vts!r \in A$ 
⟨proof⟩

lemma construct-pocket-0-last-in-set:
assumes  $\text{set}(\text{tl } vts) \cap A \neq \{\}$ 
assumes  $vts!0 \in A$ 
defines  $p \equiv \text{construct-pocket-0 } vts A$ 
shows  $\text{last } p \in A$ 
⟨proof⟩

lemma construct-pocket-0-first-last-distinct:
assumes  $\text{card } A \geq 2$ 
assumes  $A \subseteq \text{set } vts$ 
assumes  $\text{distinct}(\text{butlast } vts)$ 
assumes  $\text{hd } vts = \text{last } vts$ 
shows  $\text{hd}(\text{construct-pocket-0 } vts A) \neq \text{last}(\text{construct-pocket-0 } vts A)$ 
⟨proof⟩

lemma construct-pocket-is-pocket:
assumes  $\text{polygon}(\text{make-polygonal-path } vts)$ 
assumes  $vts!0 \in \text{frontier}(\text{convex hull}(\text{set } vts))$ 
assumes  $vts!1 \notin \text{frontier}(\text{convex hull}(\text{set } vts))$ 
shows  $\text{is-pocket-0 } vts (\text{construct-pocket-0 } vts (\text{set } vts \cap \text{frontier}(\text{convex hull}(\text{set } vts))))$ 
⟨proof⟩

lemma exists-point-above-interior:
fixes  $a :: \text{real}^2$ 
assumes  $a \in \text{interior}(\text{convex hull } S)$ 

```

**obtains**  $x$  **where**  $x \in S \wedge x\$2 > a\$2$   
 $\langle proof \rangle$

**lemma** *exists-point-above-convex-hull-interior*:  
**fixes**  $S :: (real^2)$  set  
**assumes**  $S \neq \{\}$   
**assumes** *compact*  $S$   
**obtains**  $x$  **where**  $x \in S - (\text{interior}(\text{convex hull } S)) \wedge (\forall y \in \text{interior}(\text{convex hull } S). x\$2 > y\$2)$   
 $\langle proof \rangle$

**lemma** *flip-function*:  
**defines**  $M \equiv (\text{vector}[\text{vector}[1, 0], \text{vector}[0, -1]] :: (real^2)^2)$   
**defines**  $f \equiv \lambda v. M * v v$   
**defines**  $g \equiv (\lambda v. \text{vector}[v\$1, -v\$2]) :: (real^2 \Rightarrow real^2)$   
**shows** *inj*  $f f = g$   
 $\langle proof \rangle$

**lemma** *exists-point-below-convex-hull-interior*:  
**fixes**  $S :: (real^2)$  set  
**assumes**  $S \neq \{\}$   
**assumes** *compact*  $S$   
**obtains**  $x$  **where**  $x \in S - (\text{interior}(\text{convex hull } S)) \wedge (\forall y \in \text{interior}(\text{convex hull } S). x\$2 < y\$2)$   
 $\langle proof \rangle$

**lemma** *exists-point-above-all*:  
**fixes**  $p q :: R\text{-to-}R2$   
**defines**  $H \equiv \text{convex hull}(\text{path-image } p \cup \text{path-image } q)$   
**assumes** *path*  $p \wedge \text{path } q$   
**assumes**  $p^{\{0 <.. < 1\}} \subseteq \text{interior } H$   
**assumes**  $(p 0)\$2 = 0 \wedge (p 1)\$2 = 0$   
**assumes**  $\exists x \in p^{\{0 <.. < 1\}}. x\$2 \geq 0$   
**obtains**  $x$  **where**  $x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 > y\$2)$   
 $\langle proof \rangle$

**lemma** *exists-point-below-all*:  
**fixes**  $p q :: R\text{-to-}R2$   
**defines**  $H \equiv \text{convex hull}(\text{path-image } p \cup \text{path-image } q)$   
**assumes** *path*  $p \wedge \text{path } q$   
**assumes**  $p^{\{0 <.. < 1\}} \subseteq \text{interior } H$   
**assumes**  $(p 0)\$2 = 0 \wedge (p 1)\$2 = 0$   
**assumes**  $\exists x \in \text{path-image } p \cup \text{path-image } q. x\$2 < 0$   
**obtains**  $x$  **where**  $x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 < y\$2)$   
 $\langle proof \rangle$

**lemma** *pocket-fill-line-int-aux*:  
**fixes**  $x y z :: real^2$   
**defines**  $a \equiv y\$1$

```

assumes  $x = 0$ 
assumes  $a > 0 \wedge y\$2 = 0$ 
assumes  $z\$1 < 0 \vee z\$1 > a$ 
assumes  $z\$2 = 0$ 
assumes  $\text{convex } A \wedge \text{compact } A$ 
assumes  $\{x, y, z\} \subseteq A$ 
assumes  $\{x, y\} \subseteq \text{frontier } A$ 
shows  $z \in \text{frontier } A \wedge \text{closed-segment } x y \subseteq \text{frontier } A$ 
⟨proof⟩

lemma axis-dist:
fixes  $a b :: \text{real}^2$ 
shows  $a\$2 = b\$2 \implies \text{dist } a b = \text{dist } (a\$1) (b\$1)$ 
 $a\$1 = b\$1 \implies \text{dist } a b = \text{dist } (a\$2) (b\$2)$ 
⟨proof⟩

lemma dist-bound-1:
fixes  $a b x :: \text{real}^2$ 
assumes  $a\$2 = x\$2$ 
assumes  $b \in \text{ball } x \varepsilon$ 
assumes  $\varepsilon < \text{dist } a x$ 
shows  $a\$1 < x\$1 \implies b\$1 > a\$1$ 
 $a\$1 > x\$1 \implies b\$1 < a\$1$ 
⟨proof⟩

lemma dist-bound-2:
fixes  $a b x :: \text{real}^2$ 
assumes  $a\$1 = x\$1$ 
assumes  $b \in \text{ball } x \varepsilon$ 
assumes  $\varepsilon < \text{dist } a x$ 
shows  $a\$2 < x\$2 \implies b\$2 > a\$2$ 
 $a\$2 > x\$2 \implies b\$2 < a\$2$ 
⟨proof⟩

lemma linepath-bound-1:
fixes  $x y :: \text{real}^2$ 
shows  $a < x\$1 \wedge a < y\$1 \implies \forall v \in \text{path-image } (\text{linepath } x y). a < v\$1$ 
 $x\$1 < b \wedge y\$1 < b \implies \forall v \in \text{path-image } (\text{linepath } x y). v\$1 < b$ 
⟨proof⟩

lemma linepath-bound-2:
fixes  $x y :: \text{real}^2$ 
shows  $a < x\$2 \wedge a < y\$2 \implies \forall v \in \text{path-image } (\text{linepath } x y). a < v\$2$ 
 $x\$2 < b \wedge y\$2 < b \implies \forall v \in \text{path-image } (\text{linepath } x y). v\$2 < b$ 
⟨proof⟩

lemma linepath-int-corner:
fixes  $x y z :: \text{real}^2$ 
assumes  $x\$2 \neq y\$2$ 
assumes  $y\$2 = z\$2$ 
shows  $\text{path-image } (\text{linepath } x y) \cap \text{path-image } (\text{linepath } y z) = \{y\}$ 

```

```

(is path-image ?l1 ∩ path-image ?l2 = {y})
⟨proof⟩

lemma linepath-int-vertical:
fixes w x y z :: real^2
assumes w$1 ≠ y$1
assumes w$1 = x$1
assumes y$1 = z$1
shows path-image (linepath w x) ∩ path-image (linepath y z) = {}
⟨proof⟩

lemma linepath-int-horizontal:
fixes w x y z :: real^2
assumes w$2 ≠ y$2
assumes w$2 = x$2
assumes y$2 = z$2
shows path-image (linepath w x) ∩ path-image (linepath y z) = {}
⟨proof⟩

lemma linepath-int-columns:
fixes w x y z :: real^2
assumes w$1 < y$1 ∧ w$1 < z$1
assumes x$1 < y$1 ∧ x$1 < z$1
shows path-image (linepath w x) ∩ path-image (linepath y z) = {}
(is path-image ?l1 ∩ path-image ?l2 = {})
⟨proof⟩

lemma linepath-int-rows:
fixes w x y z :: real^2
assumes w$2 < y$2 ∧ w$2 < z$2
assumes x$2 < y$2 ∧ x$2 < z$2
shows path-image (linepath w x) ∩ path-image (linepath y z) = {}
(is path-image ?l1 ∩ path-image ?l2 = {})
⟨proof⟩

lemma horizontal-segment-at-0:
assumes a > 0
shows closed-segment ((vector [0, 0])::(real^2)) (vector [a, 0]) = {x. x$2 = 0
∧ x$1 ∈ {0..a}}
(is ?l = ?s)
⟨proof⟩

lemma horizontal-segment-at-0':
fixes x y :: real^2
assumes a > 0
assumes x$1 = 0 ∧ x$2 = 0 ∧ y$1 = a ∧ y$2 = 0
shows closed-segment x y = {x. x$2 = 0 ∧ x$1 ∈ {0..a}}
⟨proof⟩

```

```

lemma pocket-fill-line-int-aux1:
  fixes p q :: R-to-R2
  defines p0 ≡ pathstart p
  defines p1 ≡ pathfinish p
  defines q0 ≡ pathstart q
  defines q1 ≡ pathfinish q
  defines a ≡ p1$1
  defines l ≡ closed-segment p0 p1
  assumes simple-path p
  assumes simple-path q
  assumes p0$1 = 0 ∧ p0$2 = 0 ∧ p1$2 = 0
  assumes a > 0
  assumes path-image q ∩ {x. x$2 = 0} ⊆ l
  assumes path-image p ∩ {x. x$2 = 0} ⊆ l
  assumes ∀ v ∈ path-image p. q0$2 ≤ v$2
  assumes ∀ v ∈ path-image p. q1$2 > v$2
  shows path-image p ∩ path-image q ≠ {}
  ⟨proof⟩

lemma pocket-fill-line-int-aux2:
  fixes p q :: R-to-R2
  fixes A :: (real2) set
  defines p0 ≡ pathstart p
  defines p1 ≡ pathfinish p
  defines a ≡ p1$1
  defines l ≡ closed-segment p0 p1
  assumes simple-path p
  assumes p0$1 = 0 ∧ p0$2 = 0 ∧ p1$2 = 0
  assumes a > 0
  assumes convex A ∧ compact A
  assumes {p0, p1} ⊆ frontier A
  assumes p ‘ {0 <..< 1} ⊆ interior A
  shows path-image p ∩ {x. x$2 = 0} ⊆ l
  ⟨proof⟩

lemma three-points-on-line:
  fixes a b :: 'a::real-vector
  assumes A = affine hull {a, b}
  assumes a ≠ b
  assumes {x, y, z} ⊆ A
  assumes x ≠ y ∧ y ≠ z ∧ x ≠ z
  shows x ∈ open-segment y z ∨ y ∈ open-segment x z ∨ z ∈ open-segment x y
  ⟨proof⟩

lemma pocket-fill-line-int-aux3:
  fixes A :: (real2) set
  assumes convex A ∧ compact A
  assumes v ≠ 0
  assumes closed-segment 0 w ⊆ frontier A (is closed-segment ?a ?b ⊆ -)

```

```

assumes  $w \cdot v = 0$ 
assumes  $w \neq 0$ 
shows ( $A \subseteq \{x. x \cdot v \leq 0\} \vee A \subseteq \{x. x \cdot v \geq 0\}$ ) (is  $A \subseteq ?P1 \vee A \subseteq ?P2$ )
⟨proof⟩

lemma pocket-fill-line-int-aux4:
  fixes  $p q :: R\text{-to-}R2$ 
  fixes  $A :: (\text{real}^2) \text{ set}$ 
  defines  $p0 \equiv \text{pathstart } p$ 
  defines  $p1 \equiv \text{pathfinish } p$ 
  defines  $q0 \equiv \text{pathstart } q$ 
  defines  $q1 \equiv \text{pathfinish } q$ 
  defines  $a \equiv p1\$1$ 
  defines  $l \equiv \text{closed-segment } p0\ p1$ 
  assumes simple-path  $p$ 
  assumes simple-path  $q$ 
  assumes path-image  $p \cap \text{path-image } q = \{\}$ 
  assumes  $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$ 
  assumes  $a > 0$ 
  assumes  $\forall v \in \text{path-image } p. q0\$2 \leq v\$2$ 
  assumes  $\forall v \in \text{path-image } p. q1\$2 > v\$2$ 
  assumes convex  $A \wedge \text{compact } A$ 
  assumes  $\{p0, p1\} \subseteq \text{frontier } A$ 
  assumes  $p^{\{0<..<1\}} \subseteq \text{interior } A$ 
  assumes path-image  $q \subseteq A$ 
  shows  $l \subseteq \text{frontier } A \forall x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 \geq 0 \quad q0\$2 = 0$ 
⟨proof⟩

lemma pocket-fill-line-int-aux5:
  fixes  $p q :: R\text{-to-}R2$ 
  fixes  $A :: (\text{real}^2) \text{ set}$ 
  defines  $p0 \equiv \text{pathstart } p$ 
  defines  $p1 \equiv \text{pathfinish } p$ 
  defines  $q0 \equiv \text{pathstart } q$ 
  defines  $q1 \equiv \text{pathfinish } q$ 
  defines  $a \equiv p1\$1$ 
  defines  $l \equiv \text{closed-segment } p0\ p1$ 
  assumes simple-path  $p$ 
  assumes simple-path  $q$ 
  assumes path-image  $p \cap \text{path-image } q = \{q0, q1\}$ 
  assumes  $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$ 
  assumes  $a > 0$ 
  assumes  $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$ 
  assumes  $\{p0, p1\} \subseteq \text{frontier } A$ 
  assumes  $p^{\{0<..<1\}} \subseteq \text{interior } A$ 
  assumes path-image  $q \subseteq A$ 
  assumes  $\exists x \in p^{\{0<..<1\}}. x\$2 \geq 0$ 
  assumes  $q0 = p1 \wedge q1 = p0$ 

```

**shows**  $l \subseteq \text{frontier } A \ \forall x \in \text{path-image } p \cup \text{path-image } q. \ x\$2 \geq 0$   
 $\langle \text{proof} \rangle$

```
lemma pocket-fill-line-int-aux6:
  fixes p q :: R-to-R2
  defines p0 ≡ pathstart p
  defines p1 ≡ pathfinish p
  defines q0 ≡ pathstart q
  defines q1 ≡ pathfinish q
  defines a ≡ p1$1
  assumes simple-path p
  assumes simple-path q
  assumes p0 = 0 ∧ p1$2 = 0
  assumes a > 0
  assumes q0$1 ∈ {0..a} ∧ q0$2 = 0
  assumes ∀x ∈ path-image p. q1$2 > x$2
  assumes ∀x ∈ path-image p ∪ path-image q. x$2 ≥ 0
  shows path-image p ∩ path-image q ≠ {}
 $\langle \text{proof} \rangle$ 
```

```
lemma pocket-fill-line-int-aux7:
  fixes p q :: R-to-R2
  fixes A :: (real^2) set
  defines p0 ≡ pathstart p
  defines p1 ≡ pathfinish p
  defines q0 ≡ pathstart q
  defines q1 ≡ pathfinish q
  defines a ≡ p1$1
  defines l ≡ open-segment p0 p1
  assumes simple-path p
  assumes simple-path q
  assumes path-image p ∩ path-image q = {q0, q1}
  assumes p0$1 = 0 ∧ p0$2 = 0 ∧ p1$2 = 0
  assumes a > 0
  assumes A = convex hull (path-image p ∪ path-image q)
  assumes {p0, p1} ⊆ frontier A
  assumes p‘{0<..<1} ⊆ interior A
  assumes ∃x ∈ p‘{0<..<1}. x$2 ≥ 0
  assumes q0 = p1 ∧ q1 = p0
  shows path-image q ∩ l = {} closed-segment p0 p1 ⊆ frontier A
 $\langle \text{proof} \rangle$ 
```

```
lemma frontier-injective-linear-image:
  fixes f :: 'a::euclidean-space ⇒ 'a::euclidean-space
  assumes linear f inj f
  shows f ‘ (frontier S) = frontier (f ‘ S)
 $\langle \text{proof} \rangle$ 
```

```

lemma pocket-fill-line-int-aux8:
  fixes p q :: R-to-R2
  fixes A :: (real2) set
  defines p0 ≡ pathstart p
  defines p1 ≡ pathfinish p
  defines q0 ≡ pathstart q
  defines q1 ≡ pathfinish q
  defines a ≡ p1$1
  defines l ≡ open-segment p0 p1
  assumes simple-path p
  assumes simple-path q
  assumes path-image p ∩ path-image q = {q0, q1}
  assumes p0$1 = 0 ∧ p0$2 = 0 ∧ p1$2 = 0
  assumes a > 0
  assumes A = convex hull (path-image p ∪ path-image q)
  assumes {p0, p1} ⊆ frontier A
  assumes p'{0<..<1} ⊆ interior A
  assumes q0 = p1 ∧ q1 = p0
  shows path-image q ∩ l = {} ∧ l ⊆ frontier A
⟨proof⟩

lemma simple-path-linear-image:
  assumes simple-path p
  assumes inj f ∧ bounded-linear f
  shows simple-path (f ∘ p)
⟨proof⟩

lemma vts-interior:
  fixes vts
  defines p ≡ make-polygonal-path vts
  assumes convex H
  assumes ∀j ∈ {0<..<length vts − 1}. vts!j ∉ frontier H
  assumes loop-free p
  assumes path-image p ⊆ H
  assumes length vts ≥ 3
  shows p'{0<..<1} ⊆ interior H
⟨proof⟩

lemma pocket-fill-line-int-0:
  assumes polygon-of r vts
  defines H ≡ convex hull (set vts)
  assumes 2 ≤ i ∧ i < length vts − 1
  defines a ≡ hd vts
  defines b ≡ vts!i
  assumes {a, b} ⊆ frontier H
  assumes ∀j ∈ {0<..<i}. vts!j ∉ frontier H
  assumes a = 0
  shows path-image (linepath a b) ∩ path-image r = {a, b}
    path-image (linepath a b) ⊆ frontier H

```

$\langle proof \rangle$

**lemma** *linepath-translation*:  $(\lambda v. v - a) \circ (\text{linepath } x y) = \text{linepath } ((\lambda v. v - a) x) ((\lambda v. v - a) y)$   
 $\langle proof \rangle$

**lemma** *linepath-image-translation*:  
 $\text{path-image } ((\lambda v. v - a) \circ (\text{linepath } x y)) = \text{path-image } (\text{linepath } ((\lambda v. v - a) x) ((\lambda v. v - a) y))$   
 $\langle proof \rangle$

**lemma** *make-polygonal-path-translate*:  
**assumes**  $\text{length } vts \geq 1$   
**shows**  $(\lambda v. v - a) \circ (\text{make-polygonal-path } vts) = \text{make-polygonal-path } (\text{map } (\lambda v. v - a) vts)$   
 $\langle proof \rangle$

**lemma** *pocket-fill-line-int*:  
**assumes** *polygon-of*  $r$   $vts$   
**defines**  $H \equiv \text{convex hull } (\text{set } vts)$   
**assumes**  $2 \leq i \wedge i < \text{length } vts - 1$   
**defines**  $a \equiv \text{hd } vts$   
**defines**  $b \equiv vts!i$   
**assumes**  $\{a, b\} \subseteq \text{frontier } H$   
**assumes**  $\forall j \in \{0 <.. < i\}. vts!j \notin \text{frontier } H$   
**shows**  $\text{path-image } (\text{linepath } a b) \cap \text{path-image } r = \{a, b\}$   
 $\text{path-image } (\text{linepath } a b) \subseteq \text{frontier } H$   
 $\langle proof \rangle$

**lemma** *path-connected-simple-path-endless*:  
**assumes** *simple-path*  $p$   
**shows** *path-connected* ( $\text{path-image } p - \{\text{pathstart } p, \text{pathfinish } p\}$ ) (**is** *path-connected*  $?S$ )  
 $\langle proof \rangle$

**lemma** *simple-loop-split*:  
**assumes** *simple-path*  $p \wedge \text{closed-path } p$   
**assumes** *simple-path*  $q$   
**assumes**  $\text{path-image } q \cap \text{path-image } p = \{q 0, q 1\}$   
**assumes**  $\text{path-image } q \cap \text{path-inside } p \neq \{\}$   
**shows**  $q^{\{0 <.. < 1\}} \subseteq \text{path-inside } p$   
 $\langle proof \rangle$

**lemma** *pocket-path-interior-aux*:  
**assumes** *simple-path*  $p \wedge \text{simple-path } q$   
**assumes** *arc*  $p \wedge \text{arc } q$   
**assumes**  $q 0 = p 1 \wedge q 1 = p 0$   
**assumes**  $\text{path-image } p \cap \text{path-image } q = \{p 0, q 0\}$

```

defines  $A \equiv \text{convex hull}(\text{path-image } p \cup \text{path-image } q)$ 
defines  $l \equiv \text{linepath}(p\ 0)\ (p\ 1)$ 
assumes  $p^{\{0 <..<1\}} \subseteq \text{interior } A$ 
assumes  $\text{path-image } l \subseteq \text{frontier } A$ 
assumes  $\text{path-image } q \cap \text{path-image } l = \{l\ 0, q\ 0\}$ 
shows  $p^{\{0 <..<1\}} \cap \text{path-inside}(l+++q) \neq \{\}$ 
    simple-path  $(l+++q) \wedge \text{closed-path} (l+++q)$ 
     $\text{path-image } p \cap \text{path-image } (l+++q) = \{p\ 0, p\ 1\}$ 
⟨proof⟩

lemma pocket-path-interior:
assumes simple-path  $p \wedge \text{simple-path } q$ 
assumes arc  $p \wedge \text{arc } q$ 
assumes  $q\ 0 = p\ 1 \wedge q\ 1 = p\ 0$ 
assumes  $\text{path-image } p \cap \text{path-image } q = \{p\ 0, q\ 0\}$ 
defines  $A \equiv \text{convex hull}(\text{path-image } p \cup \text{path-image } q)$ 
defines  $l \equiv \text{linepath}(p\ 0)\ (p\ 1)$ 
assumes  $p^{\{0 <..<1\}} \subseteq \text{interior } A$ 
assumes  $\text{path-image } l \subseteq \text{frontier } A$ 
assumes  $\text{path-image } q \cap \text{path-image } l = \{l\ 0, q\ 0\}$ 
shows  $p^{\{0 <..<1\}} \subseteq \text{path-inside}(l+++q)$ 
⟨proof⟩

lemma pocket-path-good:
assumes polygon (make-polygonal-path vts)
assumes vts!0 ∈ frontier (convex hull (set vts))
assumes vts!1 ∉ frontier (convex hull (set vts))
assumes ¬ convex (path-image (make-polygonal-path vts) ∪ path-inside (make-polygonal-path vts))
defines pocket-path-vts ≡ construct-pocket-0 vts (set vts ∩ frontier (convex hull (set vts)))
defines pocket ≡ make-polygonal-path (pocket-path-vts @ [pocket-path-vts!0])
defines filled-vts ≡ fill-pocket-0 vts (length pocket-path-vts)
defines filled-p ≡ make-polygonal-path filled-vts
defines a ≡ hd pocket-path-vts
defines b ≡ last pocket-path-vts
defines good-pocket-path-vts ≡ tl (butlast pocket-path-vts)
shows polygon filled-p
    is-polygon-split-path (butlast filled-vts) 0 1 good-pocket-path-vts
    polygon pocket
    card (set pocket-path-vts) < card (set vts)
    card (set filled-vts) < card (set vts)
⟨proof⟩

```

### 29.3 Arbitrary Polygon Case

```

lemma pick-rotate:
assumes polygon-of  $p$  vts
assumes all-integral vts

```

```

obtains  $p' vts'$  where  $\text{polygon-of } p' \text{ } vts'$ 
   $\wedge vts!0 \in \text{frontier}(\text{convex hull}(\text{set } vts'))$ 
   $\wedge \text{path-image } p' = \text{path-image } p$ 
   $\wedge \text{all-integral } vts'$ 
   $\wedge \text{set } vts' = \text{set } vts$ 
{proof}

```

**lemma** *pick-unrotated*:

```

fixes  $p :: R\text{-to-}R2$ 
assumes  $\text{polygon: polygon } p$ 
assumes  $\text{polygonal-path: } p = \text{make-polygonal-path } vts$ 
assumes  $\text{int-vertices: all-integral } vts$ 
assumes  $I\text{-is: } I = \text{card}\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ 
assumes  $B\text{-is: } B = \text{card}\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
assumes  $vts!0 \in \text{frontier}(\text{convex hull}(\text{set } vts))$ 
shows  $\text{measure lebesgue}(\text{path-inside } p) = I + B/2 - 1$ 
{proof}

```

**theorem** *pick*:

```

fixes  $p :: R\text{-to-}R2$ 
assumes  $\text{polygon } p$ 
assumes  $p = \text{make-polygonal-path } vts$ 
assumes  $\text{all-integral } vts$ 
assumes  $I = \text{card}\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ 
assumes  $B = \text{card}\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
shows  $\text{measure lebesgue}(\text{path-inside } p) = I + B/2 - 1$ 
{proof}

```

**end**

## References

- [1] B. Grünbaum and G. C. Shephard. Pick’s theorem. *The American Mathematical Monthly*, 100(2):150–161, 1993.
- [2] J. Harrison. A formal proof of Pick’s theorem. *Math. Struct. Comput. Sci.*, 21(4):715–729, 2011.