

Pick's Theorem

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Abstract

We formalize Pick's theorem for finding the area of a simple polygon whose vertices are integral lattice points [1]. We are inspired by John Harrison's formalization of Pick's theorem in HOL Light [2], but tailor our proof approach to avoid a primary challenge point in his formalization, which is proving that any polygon with more than three vertices can be split (in its interior) by a line between some two vertices. Our formalization involves augmenting the existing geometry libraries in various foundational ways (e.g., by adding the definition of a polygon and formalizing some key properties thereof).

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theory <i>Integral-Matrix</i>	
imports	
<i>Complex-Main</i>	
<i>HOL-Analysis.Finite-Cartesian-Product</i>	
<i>HOL-Analysis.Linear-Algebra</i>	
<i>HOL-Analysis.Determinants</i>	
begin	

1 Misc. Linear Algebra Setup

lemma *vec-scaleR-2*: $(c::real) *_R ((vector [a, b])::real^2) = vector [a * c, b * c]$
proof–
have $(c *_R (vector [a, b])::real^2)\$1 = a * c$ **by** *simp*
moreover have $(c *_R (vector [a, b])::real^2)\$2 = ((vector [a, b])::real^2)\$2 * c$ **by** *simp*
ultimately show *?thesis* **by** (*smt (verit, best) exhaust-2 vec-eq-iff vector-2(1) vector-2(2)*)
qed

definition *is-int* :: $real \Rightarrow bool$ **where**
is-int $x \longleftrightarrow (\exists n::int. x = n)$

lemma *is-int-sum*: $is-int\ x \wedge is-int\ y \longrightarrow is-int\ (x + y)$
by (*metis is-int-def of-int-add*)

lemma *is-int-minus*: $is-int\ x \wedge is-int\ y \longrightarrow is-int\ (x - y)$
by (*metis is-int-def of-int-diff*)

lemma *is-int-mult*: $is-int\ x \wedge is-int\ y \longrightarrow is-int\ (x * y)$
by (*metis is-int-def of-int-mult*)

definition *integral-vec* :: $real^2 \Rightarrow bool$ **where**
integral-vec $v \longleftrightarrow (is-int\ (v\$1) \wedge is-int\ (v\$2))$

lemma *integral-vec-sum*: $integral-vec\ v \wedge integral-vec\ w \longrightarrow integral-vec\ (v + w)$

proof(*rule impI*)
fix $v\ w :: real^2$
let $?x = v + w$
assume $integral-vec\ v \wedge integral-vec\ w$
then obtain $v1\ v2\ w1\ w2 :: int$ **where** $v\$1 = v1 \wedge v\$2 = v2 \wedge w\$1 = w1 \wedge w\$2 = w2$
using *integral-vec-def is-int-def* **by** *auto*
then have $?x\$1 = v1 + w1$ **and** $?x\$2 = v2 + w2$ **by** *auto*
thus $integral-vec\ ?x$ **using** *integral-vec-def is-int-def* **by** *blast*
qed

lemma *integral-vec-minus*: $integral-vec\ v \longrightarrow integral-vec\ (-v)$

proof(*rule impI*)
assume $integral-vec\ v$
then obtain $x\ y :: int$ **where** $v\$1 = x \wedge v\$2 = y$
using *integral-vec-def is-int-def* **by** *auto*
then have $(-v)\$1 = -x$ **and** $(-v)\$2 = -y$
using *integral-vec-def is-int-def* **by** *auto*
thus $integral-vec\ (-v)$
using *integral-vec-def is-int-def* **by** *blast*
qed

lemma *real-2-inner*:

shows $((\text{vector } [a, b]) :: (\text{real}^2)) \cdot ((\text{vector } [c, d]) :: (\text{real}^2)) = a*c + b*d$
(is $?v \cdot ?w = a*c + b*d$)

proof –

have $?v \cdot ?w = (\sum i \in \text{UNIV}. ?v\$i \cdot ?w\$i)$ **using** *inner-vec-def*[of $?v ?w$] **by**
blast

moreover have $\forall i. ?v\$i \cdot ?w\$i = ?v\$i * ?w\i **using** *inner-real-def* **by** *simp*

ultimately have $?v \cdot ?w = (\sum i \in \text{UNIV}. ?v\$i * ?w\$i)$ **by** *presburger*

thus *thesis* **by** (*simp add: sum-2*)

qed

lemma *integral-vec-2*:

fixes $a b :: \text{int}$

assumes $v = \text{vector } [a, b]$

shows *integral-vec* v

by (*simp add: assms is-int-def integral-vec-def*)

definition *matrix-inv* $:: \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{bool}$ **where**

matrix-inv $A A' \longleftrightarrow (A ** A' = \text{mat } 1 \wedge A' ** A = \text{mat } 1)$

lemma *mat-vec-mult-2*:

fixes $v :: \text{real}^2$ **and**

$T :: \text{real}^2 \Rightarrow \text{real}^2$

defines $x: x \equiv v\$1$ **and** $y: y \equiv v\$2$ **and**

$a: a \equiv T\$1\1 **and** $b: b \equiv T\$1\2 **and**

$c: c \equiv T\$2\1 **and** $d: d \equiv T\$2\2

shows $(T *v v) = \text{vector } [x*a + y*b, x*c + y*d]$

proof –

have $(T *v v)\$1 = x*a + y*b$ **by** (*simp add: a b matrix-vector-mult-def sum-2*
 $x y$)

moreover have $(T *v v)\$2 = x*c + y*d$ **by** (*simp add: c d matrix-vector-mult-def*
 $sum-2 x y$)

ultimately show $T *v v = \text{vector } [x*a + y*b, x*c + y*d]$

by (*smt (verit) exhaust-2 vec-eq-iff vector-2(1) vector-2(2)*)

qed

definition *integral-mat* $:: \text{real}^2 \Rightarrow \text{bool}$ **where**

integral-mat $T \longleftrightarrow (\forall v. \text{integral-vec } v \longrightarrow \text{integral-vec } (T *v v))$

definition *integral-mat-surj* $:: \text{real}^2 \Rightarrow \text{bool}$ **where**

integral-mat-surj $T \longleftrightarrow (\forall v. \text{integral-vec } v \longrightarrow (\exists w. \text{integral-vec } w \wedge T *v w = v))$

definition *integral-mat-bij* $:: \text{real}^2 \Rightarrow \text{bool}$ **where**

integral-mat-bij $T \longleftrightarrow \text{integral-mat } T \wedge \text{integral-mat-surj } T$

lemma *integral-mat-integral-vec*: $\text{integral-mat } A \longrightarrow \text{integral-vec } v \longrightarrow \text{integral-vec}$
 $(A *v v)$

using *integral-mat-def* **by** *blast*

lemma *integral-mat-int-entries*:

fixes $T :: \text{real}^{\mathbb{Z}^2}$

assumes *integral-mat* T

defines $a :: \text{int}$ $b :: \text{int}$ **and** $c :: \text{int}$ $d :: \text{int}$ **and**

$c :: \text{int}$ $d :: \text{int}$ **and** $d :: \text{int}$ $d :: \text{int}$

shows $\text{is-int } a \wedge \text{is-int } b \wedge \text{is-int } c \wedge \text{is-int } d$

proof –

let $?v = \text{vector } [1, 0]$

have *integral-vec* $(?v)$ **using** *integral-vec-2*[*of* $?v$ 1 0] **by** *auto*

then have *integral-vec* $(T *v ?v)$ **using** *assms integral-mat-def* **by** *blast*

moreover have $T *v ?v = \text{vector } [a, c]$

using *mat-vec-mult-2*[*of* T $?v$] a b c d **by** *auto*

ultimately have *integral-vec* $(\text{vector } [a, c])$ **by** *auto*

then have 1: $\text{is-int } a \wedge \text{is-int } c$ **using** *integral-vec-def* **by** *auto*

let $?w = \text{vector } [0, 1]$

have *integral-vec* $(?w)$ **using** *integral-vec-2*[*of* $?w$ 0 1] **by** *auto*

then have *integral-vec* $(T *v ?w)$ **using** *assms integral-mat-def* **by** *blast*

moreover have $T *v ?w = \text{vector } [b, d]$

using *mat-vec-mult-2*[*of* T $?w$] a b c d **by** *auto*

ultimately have *integral-vec* $(\text{vector } [b, d])$ **by** *auto*

then have 2: $\text{is-int } b \wedge \text{is-int } d$ **using** *integral-vec-def* **by** *auto*

thus *thesis* **using** 1 2 **by** *auto*

qed

2 Integral Bijective Matrix Determinant

lemma *integral-mat-int-det*:

fixes $T :: \text{real}^{\mathbb{Z}^2}$

assumes *integral-mat* T

shows $\text{is-int } (\text{det } T)$

proof –

obtain a b c d **where** $abcd$: $T\$1\$1 = a \wedge T\$1\$2 = b \wedge T\$2\$1 = c \wedge T\$2\$2 = d$ **by** *auto*

have *abcd-int*: $\text{is-int } a \wedge \text{is-int } b \wedge \text{is-int } c \wedge \text{is-int } d$

using *integral-mat-int-entries*[*of* T] $abcd$ *assms* **by** *auto*

obtain ai bi ci di $:: \text{int}$ **where** $abcdi$: $ai = a \wedge bi = b \wedge ci = c \wedge di = d$

using *abcd-int is-int-def* **by** *auto*

have $\text{det } T = a*d - b*c$ **using** *det-2*[*of* T] $abcd$ **by** *auto*

also have $\dots = ai*di - bi*ci$ **using** $abcdi$ **by** *auto*

finally show *thesis* **using** *is-int-def* **by** *blast*

qed

lemma *integral-mat-bij-inv*:

fixes $T :: \text{real}^{\mathbb{Z}^2}$

assumes *integral-mat-bij* T

obtains $Tinv$ **where** $invertible\ T \wedge integral\text{-}mat\text{-}bij\ Tinv \wedge matrix\text{-}inv\ T\ Tinv$
proof –

let $?e1 = vector\ [1, 0]$
let $?e2 = vector\ [0, 1]$
let $?I = (vector\ [?e1, ?e2])::(real^{2^2})$
have $id: ?I = ((mat\ 1)::(real^{2^2}))$
unfolding $vec\text{-}eq\text{-}iff$
by $(smt\ (verit, ccfv\text{-}threshold)\ exhaust\text{-}2\ mat\text{-}def\ vec\text{-}lambda\text{-}beta\ vector\text{-}2)$
have $integral\text{-}vec\ ?e1$
by $(simp\ add: integral\text{-}vec\text{-}def\ is\text{-}int\text{-}def)$
moreover **have** $integral\text{-}vec\ ?e2$
by $(simp\ add: integral\text{-}vec\text{-}def\ is\text{-}int\text{-}def)$
ultimately obtain $x\ y$ **where** $xy: T *v\ x = ?e1 \wedge integral\text{-}vec\ x \wedge T *v\ y = ?e2 \wedge integral\text{-}vec\ y$
by $(meson\ assms\ integral\text{-}mat\text{-}bij\text{-}def\ integral\text{-}mat\text{-}surj\text{-}def)$

let $?Tinv = transpose\ (vector\ [x, y])::(real^{2^2})$
have $T ** ?Tinv = mat\ 1$ (**is** $?TxTinv = mat\ 1$)

proof –

have $column\ 1\ ?TxTinv = T *v\ (column\ 1\ ?Tinv)$
by $(metis\ matrix\text{-}vector\text{-}mul\text{-}assoc\ matrix\text{-}vector\text{-}mult\text{-}basis)$
also **have** $\dots = T *v\ x$
by $(simp\ add: row\text{-}def)$
finally **have** $[simp]: column\ 1\ ?TxTinv = ?e1$
using xy **by** $presburger$

have $column\ 2\ ?TxTinv = T *v\ (column\ 2\ ?Tinv)$
by $(metis\ matrix\text{-}vector\text{-}mul\text{-}assoc\ matrix\text{-}vector\text{-}mult\text{-}basis)$
also **have** $\dots = T *v\ y$
by $(simp\ add: row\text{-}def)$
finally **have** $[simp]: column\ 2\ ?TxTinv = ?e2$
using xy **by** $presburger$

have $\forall v. ?TxTinv *v\ v = v$

proof $(rule\ allI)$

fix $v :: real^{2^2}$

have $(?TxTinv *v\ v)\$1 = (column\ 1\ ?TxTinv)\$1 * v\$1 + (column\ 2\ ?TxTinv)\$1 * v\$2$

by $(metis\ (no\text{-}types, lifting)\ cart\text{-}eq\text{-}inner\text{-}axis\ mat\text{-}vec\text{-}mult\text{-}2\ matrix\text{-}vector\text{-}mul\text{-}component\ matrix\text{-}vector\text{-}mult\text{-}basis\ mult.\ commute\ vector\text{-}2(1))$

also **have** $\dots = v\$1$ **by** $simp$

finally **have** $v1: (?TxTinv *v\ v)\$1 = v\1 .

have $(?TxTinv *v\ v)\$2 = (column\ 1\ ?TxTinv)\$2 * v\$1 + (column\ 2\ ?TxTinv)\$2 * v\$2$

by $(metis\ (no\text{-}types, lifting)\ cart\text{-}eq\text{-}inner\text{-}axis\ mat\text{-}vec\text{-}mult\text{-}2\ matrix\text{-}vector\text{-}mul\text{-}component\ matrix\text{-}vector\text{-}mult\text{-}basis\ mult.\ commute\ vector\text{-}2(2))$

also **have** $\dots = v\$2$ **by** $simp$

finally have $v2: (?TxTinv * v) \$2 = v \2 .

show $?TxTinv * v v = v$ **using** $v1 v2$ **by** (*metis mat-vec-mult-2 matrix-vector-mul-lid*)

qed

thus $?thesis$ **by** (*simp add: matrix-eq*)

qed

then have *matrix-inv T ?Tinv*

by (*simp add: Integral-Matrix.matrix-inv-def matrix-left-right-inverse*)

moreover have *invertible T* **using** *calculation invertible-def matrix-inv-def* **by**

blast

moreover have *integral-mat-bij ?Tinv*

by (*smt (verit, del-insts) <T ** Finite-Cartesian-Product.transpose (vector [x, y]) = mat 1 > assms integral-mat-bij-def integral-mat-def integral-mat-surj-def matrix-left-right-inverse matrix-mul-lid matrix-vector-mul-assoc*)

ultimately show $?thesis$

using $<T ** Finite-Cartesian-Product.transpose (vector [x, y]) = mat 1 >$ *invertible-right-inverse* **that** **by** *blast*

qed

lemma *integral-mat-bij-det-pm1*:

fixes $T :: real^{2^2}$

assumes *integral-mat-bij T*

shows $det T = 1 \vee det T = -1$

proof–

obtain $Tinv$ **where** $Tinv: invertible T \wedge integral-mat-bij Tinv \wedge matrix-inv T Tinv$

using *integral-mat-bij-inv[of T] assms* **by** *auto*

moreover have *is-int (det Tinv)*

using *integral-mat-bij-def integral-mat-int-det[of Tinv] calculation* **by** *auto*

moreover have *is-int (det T)*

using *integral-mat-bij-def integral-mat-int-det[of T] assms* **by** *auto*

moreover have $det Tinv = 1 / det T$

proof–

have $id: Tinv ** T = mat 1$ **using** $Tinv$ **unfolding** *matrix-inv-def invertible-def*

by (*simp add: verit-sko-ex'*)

have $det Tinv * det T = det (Tinv ** T)$ **by** (*simp add: det-mul*)

also have $\dots = det ((mat 1)::real^{2^2})$ **using** id **by** *auto*

also have $\dots = (1::real)$ **by** *auto*

finally have $det Tinv * det T = 1$.

thus $?thesis$ **using** *invertible-det-nz nonzero-eq-divide-eq* **by** *fastforce*

qed

ultimately have $T-Tinv-int: is-int (det T) \wedge is-int (1 / det T)$ **by** *auto*

thus $det T = 1 \vee det T = -1$

proof–

have $abs (det T) \leq 1$ (**is** $?D \leq 1$)

proof(*rule ccontr*)

assume $\neg ?D \leq 1$

then have $?D > 1$ **by** *auto*

```

moreover from this have  $1 / ?D < 1$  by auto
moreover from calculation have  $1 / ?D > 0$  by auto
ultimately have  $\neg$  is-int  $(1 / ?D)$  unfolding is-int-def by force
moreover from T-Inv-int have is-int  $(1 / ?D)$ 
  by (smt (verit)  $\langle 1 / |\det T| < 1 \rangle$  abs-div-pos abs-divide abs-ge-self
abs-minus-cancel divide-cancel-left divide-pos-neg int-less-real-le is-int-def of-int-code(2))
  ultimately show False by auto
qed
then have  $\det T \geq -1 \wedge \det T \leq 1$ 
  using assms by auto
moreover have  $\det T \neq 0$  using integral-mat-bij-inv invertible-det-nz assms
by auto
  ultimately show  $\det T = 1 \vee \det T = -1$  using is-int-def T-Inv-int by
auto
  qed
qed

end
theory Polygon-Jordan-Curve
imports
  HOL-Analysis.Cartesian-Space
  HOL-Analysis.Path-Connected
  Poincare-Bendixson.Poincare-Bendixson
  Integral-Matrix

```

```

begin

```

3 Polygon Definitions

```

type-synonym R-to-R2 =  $(\text{real} \Rightarrow \text{real}^2)$ 

```

```

definition closed-path :: R-to-R2  $\Rightarrow$  bool where
  closed-path g  $\longleftrightarrow$  path g  $\wedge$  pathstart g = pathfinish g

```

```

definition path-inside :: R-to-R2  $\Rightarrow$   $(\text{real}^2)$  set where
  path-inside g = inside (path-image g)

```

```

definition path-outside :: R-to-R2  $\Rightarrow$   $(\text{real}^2)$  set where
  path-outside g = outside (path-image g)

```

```

fun make-polygonal-path ::  $(\text{real}^2)$  list  $\Rightarrow$  R-to-R2 where
  make-polygonal-path [] = linepath 0 0
| make-polygonal-path [a] = linepath a a
| make-polygonal-path [a,b] = linepath a b
| make-polygonal-path (a # b # xs) = (linepath a b) +++ make-polygonal-path (b
# xs)

```

```

definition polygonal-path :: R-to-R2  $\Rightarrow$  bool where
  polygonal-path g  $\longleftrightarrow$   $g \in$  make-polygonal-path{xs ::  $(\text{real}^2)$  list. True}

```

definition *all-integral* :: (real^2) list \Rightarrow bool **where**

all-integral l = $(\forall x \in \text{set } l. \text{integral-vec } x)$

definition *polygon* :: $R\text{-to-}R^2 \Rightarrow$ bool **where**

polygon g \longleftrightarrow *polygonal-path* g \wedge *simple-path* g \wedge *closed-path* g

definition *integral-polygon* :: $R\text{-to-}R^2 \Rightarrow$ bool **where**

integral-polygon g \longleftrightarrow

(*polygon* g \wedge $(\exists \text{vts}. g = \text{make-polygonal-path } \text{vts} \wedge \text{all-integral } \text{vts})$)

definition *make-triangle* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow R\text{-to-}R^2$ **where**

make-triangle a b c = *make-polygonal-path* [a, b, c, a]

definition *polygon-of* :: $R\text{-to-}R^2 \Rightarrow (\text{real}^2)$ list \Rightarrow bool **where**

polygon-of p vts \longleftrightarrow *polygon* p \wedge p = *make-polygonal-path* vts

definition *good-linepath* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow (\text{real}^2)$ list \Rightarrow bool **where**

good-linepath a b vts \longleftrightarrow (let p = *make-polygonal-path* vts in

a \neq b \wedge {a, b} \subseteq set vts \wedge *path-image* (*linepath* a b) \subseteq *path-inside* p \cup {a, b})

definition *good-polygonal-path* :: $\text{real}^2 \Rightarrow (\text{real}^2)$ list $\Rightarrow \text{real}^2 \Rightarrow (\text{real}^2)$ list \Rightarrow bool **where**

good-polygonal-path a cutvts b vts \longleftrightarrow (

let p = *make-polygonal-path* vts in

let p-cut = *make-polygonal-path* ([a] @ cutvts @ [b]) in

(a \neq b \wedge {a, b} \subseteq set vts \wedge *path-image* (p-cut) \subseteq *path-inside* p \cup {a, b} \wedge *loop-free* p-cut))

4 Jordan Curve Theorem for Polygons

definition *inside-outside* :: $R\text{-to-}R^2 \Rightarrow (\text{real}^2)$ set $\Rightarrow (\text{real}^2)$ set \Rightarrow bool **where**

inside-outside p ins outs \longleftrightarrow

(ins \neq {} \wedge *open* ins \wedge *connected* ins \wedge

outs \neq {} \wedge *open* outs \wedge *connected* outs \wedge

bounded ins \wedge \neg *bounded* outs \wedge

ins \cap outs = {} \wedge ins \cup outs = \neg *path-image* p \wedge

frontier ins = *path-image* p \wedge *frontier* outs = *path-image* p)

lemma *Jordan-inside-outside-real2*:

fixes p :: $\text{real} \Rightarrow \text{real}^2$

assumes *simple-path* p *pathfinish* p = *pathstart* p

shows *inside*(*path-image* p) \neq {} \wedge

open(*inside*(*path-image* p)) \wedge

connected(*inside*(*path-image* p)) \wedge

outside(*path-image* p) \neq {} \wedge

open(*outside*(*path-image* p)) \wedge

connected(*outside*(*path-image* p)) \wedge

$$\begin{aligned}
& \text{bounded}(\text{inside}(\text{path-image } p)) \wedge \\
& \neg \text{bounded}(\text{outside}(\text{path-image } p)) \wedge \\
& \text{inside}(\text{path-image } p) \cap \text{outside}(\text{path-image } p) = \{\} \wedge \\
& \text{inside}(\text{path-image } p) \cup \text{outside}(\text{path-image } p) = \\
& \quad - \text{path-image } p \wedge \\
& \text{frontier}(\text{inside}(\text{path-image } p)) = \text{path-image } p \wedge \\
& \text{frontier}(\text{outside}(\text{path-image } p)) = \text{path-image } p
\end{aligned}$$

proof –

have *good-type*: *c1-on-open-R2-axioms* *TYPE*((*real*, 2) *vec*)

unfolding *c1-on-open-R2-axioms-def* **by** *auto*

have $\text{inside}(\text{path-image } p) \neq \{\}$ \wedge
 $\text{open}(\text{inside}(\text{path-image } p)) \wedge$
 $\text{connected}(\text{inside}(\text{path-image } p)) \wedge$
 $\text{outside}(\text{path-image } p) \neq \{\}$ \wedge
 $\text{open}(\text{outside}(\text{path-image } p)) \wedge$
 $\text{connected}(\text{outside}(\text{path-image } p)) \wedge$
 $\text{bounded}(\text{inside}(\text{path-image } p)) \wedge$
 $\neg \text{bounded}(\text{outside}(\text{path-image } p)) \wedge$
 $\text{inside}(\text{path-image } p) \cap \text{outside}(\text{path-image } p) = \{\}$ \wedge
 $\text{inside}(\text{path-image } p) \cup \text{outside}(\text{path-image } p) =$
 $\quad - \text{path-image } p \wedge$
 $\text{frontier}(\text{inside}(\text{path-image } p)) = \text{path-image } p \wedge$
 $\text{frontier}(\text{outside}(\text{path-image } p)) = \text{path-image } p$

using *assms c1-on-open-R2.Jordan-inside-outside-R2*[*of - - p*]

unfolding *c1-on-open-R2-def c1-on-open-euclidean-def c1-on-open-def* **using**

good-type

by (*metis continuous-on-empty equals0D open-empty*)

then show *?thesis unfolding inside-outside-def*

using *path-inside-def path-outside-def* **by** *auto*

qed

lemma *inside-outside-polygon*:

fixes $p :: R\text{-to-}R^2$

assumes *polygon*: *polygon* p

shows *inside-outside* p (*path-inside* p) (*path-outside* p)

proof –

have *good-type*: *c1-on-open-R2-axioms* *TYPE*((*real*, 2) *vec*)

unfolding *c1-on-open-R2-axioms-def* **by** *auto*

have *simple-path* p *pathfinish* $p = \text{pathstart } p$ **using** *assms polygon-def closed-path-def*
by *auto*

then show *?thesis using Jordan-inside-outside-real2 unfolding inside-outside-def*

using *path-inside-def path-outside-def* **by** *auto*

qed

lemma *inside-outside-unique*:

fixes $p :: R\text{-to-}R^2$

assumes *polygon* p

assumes *io1*: *inside-outside* p *inside1* *outside1*

assumes $io2$: *inside-outside* p $inside2$ $outside2$
shows $inside1 = inside2 \wedge outside1 = outside2$
proof –
have $inner1$: *inside*(*path-image* p) = $inside1$
using *dual-order.antisym* *inside-subset* *interior-eq* *interior-inside-frontier*
using $io1$ **unfolding** *inside-outside-def*
by *metis*
have $inner2$: *inside*(*path-image* p) = $inside2$
using *dual-order.antisym* *inside-subset* *interior-eq* *interior-inside-frontier*
using $io2$ **unfolding** *inside-outside-def*
by *metis*
have $eq1$: $inside1 = inside2$
using $inner1$ $inner2$
by *auto*
have $h1$: $inside1 \cup outside1 = -\text{path-image } p$
using $io1$ **unfolding** *inside-outside-def* **by** *auto*
have $h2$: $inside1 \cap outside1 = \{\}$
using $io1$ **unfolding** *inside-outside-def* **by** *auto*
have $outer1$: *outside*(*path-image* p) = $outside1$
using $io1$ $inner1$ **unfolding** *inside-outside-def*
using $h1$ $h2$ *outside-inside* **by** *auto*
have $h3$: $inside2 \cup outside2 = -\text{path-image } p$
using $io2$ **unfolding** *inside-outside-def* **by** *auto*
have $h4$: $inside2 \cap outside2 = \{\}$
using $io2$ **unfolding** *inside-outside-def* **by** *auto*
have $outer2$: *outside*(*path-image* p) = $outside2$
using $io2$ $inner2$ **unfolding** *inside-outside-def*
using $h3$ $h4$ *outside-inside* **by** *auto*
then have $eq2$: $outside1 = outside2$
using $outer1$ $outer2$ **by** *auto*
then show *?thesis* **using** $eq1$ $eq2$ **by** *auto*
qed

lemma *polygon-jordan-curve*:

fixes p :: *R-to-R2*

assumes *polygon* p

obtains *inside* *outside* **where**

inside-outside p *inside* *outside*

proof –

have *good-type*: *c1-on-open-R2-axioms* *TYPE*((*real*, 2) *vec*)

unfolding *c1-on-open-R2-axioms-def* **by** *auto*

have *simple-path* p *pathfinish* $p = \text{pathstart } p$ **using** *assms* *polygon-def* *closed-path-def*
by *auto*

then obtain *inside* *outside* **where**

inside $\neq \{\}$ *open* *inside* *connected* *inside*

outside $\neq \{\}$ *open* *outside* *connected* *outside*

bounded *inside* \neg *bounded* *outside* *inside* \cap *outside* = $\{\}$

inside \cup *outside* = $-\text{path-image } p$

frontier *inside* = *path-image* p

frontier outside = path-image p
using *c1-on-open-R2.Jordan-curve-R2*[of - - - p]
unfolding *c1-on-open-R2-def c1-on-open-euclidean-def c1-on-open-def* **using**
good-type
by (*metis continuous-on-empty equals0D open-empty*)
then show *?thesis*
using *inside-outside-def* **that by auto**
qed

lemma *connected-component-image:*

fixes $f :: 'a::euclidean-space \Rightarrow 'b::euclidean-space$

assumes *linear f bij f*

shows $f \text{ ` } (\text{connected-component-set } S \ x) = \text{connected-component-set } (f \text{ ` } S) (f \ x)$

proof –

have $\text{conn}: \bigwedge S. \text{connected } S \Longrightarrow \text{connected } (f \text{ ` } S)$

by (*simp add: assms(1) connected-linear-image*)

then have $h1: \bigwedge T. T \in \{T. \text{connected } T \wedge x \in T \wedge T \subseteq S\} \Longrightarrow f \text{ ` } T \in \{T. \text{connected } T \wedge (f \ x) \in T \wedge T \subseteq (f \text{ ` } S)\}$

by auto

then have $\text{subset1}: f \text{ ` } \text{connected-component-set } S \ x \subseteq \text{connected-component-set } (f \text{ ` } S) (f \ x)$

using *connected-component-Union*

by (*smt (verit, ccfv-threshold) assms(2) bij-is-inj connected-component-eq-empty connected-component-maximal connected-component-refl-eq connected-component-subset connected-connected-component image-is-empty inj-image-mem-iff mem-Collect-eq*)

have $\bigwedge S. \text{connected } (f \text{ ` } S) \Longrightarrow \text{connected } S$

using *assms connected-continuous-image assms linear-continuous-on linear-conv-bounded-linear bij-is-inj homeomorphism-def linear-homeomorphism-image*

by (*smt (verit, del-insts)*)

then have $h2: \bigwedge T. f \text{ ` } T \in \{T. \text{connected } T \wedge (f \ x) \in T \wedge T \subseteq (f \text{ ` } S)\} \Longrightarrow T \in \{T. \text{connected } T \wedge x \in T \wedge T \subseteq S\}$

by (*simp add: assms(2) bij-is-inj image-subset-iff inj-image-mem-iff subsetI*)

then have $\text{subset2}: \text{connected-component-set } (f \text{ ` } S) (f \ x) \subseteq f \text{ ` } \text{connected-component-set } S \ x$

using *connected-component-Union*[of $S \ x$] *connected-component-Union*[of $f \text{ ` } S \ f \ x$]

by (*smt (verit, del-insts) assms(2) bij-is-inj connected-component-eq-empty connected-component-maximal connected-component-refl-eq connected-component-subset connected-connected-component image-mono inj-image-mem-iff mem-Collect-eq subset-imageE*)

show $f \text{ ` } (\text{connected-component-set } S \ x) = \text{connected-component-set } (f \text{ ` } S) (f \ x)$

using *subset1 subset2* **by auto**

qed

lemma *bounded-map:*

fixes $f :: 'a::euclidean-space \Rightarrow 'b::euclidean-space$

assumes *linear f bij f*

shows $\text{bounded } (f \text{ ` } S) = \text{bounded } S$

```

proof –
  have h1: bounded S  $\implies$  bounded (f ' S)
    using assms
    using bounded-linear-image linear-conv-bounded-linear by blast
  have bounded-linear f
    using linear-conv-bounded-linear assms by auto
  then have bounded-linear (inv f)
    using assms unfolding bij-def
    by (smt (verit, ccfv-threshold) bij-betw-def bij-betw-subset dim-image-eq inv-equality
linear-conv-bounded-linear linear-surjective-isomorphism subset-UNIV)
  then have h2: bounded (f ' S)  $\implies$  bounded S
    using assms
    by (metis bij-is-inj bounded-linear-image image-inv-f-f)
  then show ?thesis
    using assms h1 h2 by auto
qed

```

```

lemma inside-bijective-linear-image:
  fixes f :: 'a::euclidean-space  $\Rightarrow$  'b::euclidean-space
  fixes c :: real  $\Rightarrow$  'a
  assumes c-simple:path c
  assumes linear f bij f
  shows inside (f ' (path-image c)) = f ' (inside(path-image c))
proof –
  have set1:  $\{x. x \notin f ' \text{path-image } c\} = f ' \{x. x \notin \text{path-image } c\}$ 
    using assms path-image-compose unfolding bij-def
    by (smt (verit, best) UNIV-I imageE inj-image-mem-iff mem-Collect-eq subsetI
subset-antisym)
  have linear-inv: linear (inv f)
    using assms
    by (metis bij-imp-bij-inv bij-is-inj inv-o-cancel linear-injective-left-inverse o-inv-o-cancel)
  have bij-inv: bij (inv f)
    using assms
    using bij-imp-bij-inv by blast
  have inset1:  $\bigwedge x. x \in \{x. \text{bounded (connected-component-set (- f ' path-image c) x)}\} \implies x \in f ' \{x. \text{bounded (connected-component-set (- path-image c) x)}\}$ 
proof –
  fix x
  assume  $*$ :  $x \in \{x. \text{bounded (connected-component-set (- f ' path-image c) x)}\}$ 
  have inj f
    using assms(3) bij-betw-imp-inj-on by blast
  then show  $x \in f ' \{x. \text{bounded (connected-component-set (- path-image c) x)}\}$ 
    using  $*$  connected-component-image[OF linear-inv bij-inv]
    by (smt (z3)  $\langle \bigwedge x S. \text{inv } f ' \text{connected-component-set } S \ x = \text{connected-component-set (inv } f ' S) (inv f \ x) \rangle \langle \text{bij (inv } f) \rangle \langle \text{linear (inv } f) \rangle \langle x \in \{x. \text{bounded (connected-component-set (- f ' path-image c) x)}\} \rangle \text{bij-image-Compl-eq bounded-map connected-component-eq-empty image-empty image-inv-f-f mem-Collect-eq}$ )
  qed
  have inset2:  $\bigwedge x. x \in f ' \{x. \text{bounded (connected-component-set (- path-image$ 

```

```

c) x)}  $\implies x \in \{x. \text{bounded} (\text{connected-component-set} (- f \text{' path-image } c) x)\}$ 
proof -
  fix x
  assume  $x \in f \text{' } \{x. \text{bounded} (\text{connected-component-set} (- \text{path-image } c) x)\}$ 
  then obtain x1 where  $x = f \text{ } x1$   $x1 \in \{x. \text{bounded} (\text{connected-component-set}$ 
 $(- \text{path-image } c) x)\}$ 
  by auto
  then show  $x \in \{x. \text{bounded} (\text{connected-component-set} (- f \text{' path-image } c) x)\}$ 

  using bounded-map[OF assms(2) assms(3)] connected-component-image[OF
assms(2) assms(3)]
  by (metis assms(3) bij-image-Compl-eq mem-Collect-eq)
qed
  have set2:  $f \text{' } \{x. \text{bounded} (\text{connected-component-set} (- \text{path-image } c) x)\} = \{x.$ 
 $\text{bounded} (\text{connected-component-set} (- f \text{' path-image } c) x)\}$ 
  using inset1 inset2 by auto
  have inset1:  $\bigwedge x. x \in f \text{' } \{x. x \notin \text{path-image } c \wedge \text{bounded} (\text{connected-component-set}$ 
 $(- \text{path-image } c) x)\} \implies$ 
 $x \in \{x. x \notin f \text{' path-image } c \wedge \text{bounded} (\text{connected-component-set} (- f \text{'$ 
 $\text{path-image } c) x)\}$ 
  proof -
  fix x
  assume  $x \in f \text{' } \{x. x \notin \text{path-image } c \wedge \text{bounded} (\text{connected-component-set} (-$ 
 $\text{path-image } c) x)\}$ 
  then show  $x \in \{x. x \notin f \text{' path-image } c \wedge \text{bounded} (\text{connected-component-set}$ 
 $(- f \text{' path-image } c) x)\}$ 
  by (metis (no-types, lifting) image-iff mem-Collect-eq set1 set2)
qed
  have inset2:  $\bigwedge x. x \in \{x. x \notin f \text{' path-image } c \wedge \text{bounded} (\text{connected-component-set}$ 
 $(- f \text{' path-image } c) x)\} \implies$ 
 $x \in f \text{' } \{x. x \notin \text{path-image } c \wedge \text{bounded} (\text{connected-component-set} (- \text{path-image}$ 
 $c) x)\}$ 
  proof -
  fix x
  assume  $x \in \{x. x \notin f \text{' path-image } c \wedge \text{bounded} (\text{connected-component-set} (-$ 
 $f \text{' path-image } c) x)\}$ 
  then show  $x \in f \text{' } \{x. x \notin \text{path-image } c \wedge \text{bounded} (\text{connected-component-set}$ 
 $(- \text{path-image } c) x)\}$ 
  by (smt (verit, best) image-iff mem-Collect-eq set2)
qed
  have same-set:  $\{x. x \notin f \text{' path-image } c \wedge \text{bounded} (\text{connected-component-set} (-$ 
 $f \text{' path-image } c) x)\} =$ 
 $f \text{' } \{x. x \notin \text{path-image } c \wedge \text{bounded} (\text{connected-component-set} (- \text{path-image } c)$ 
 $x)\}$ 
  using inset1 inset2
  by blast
  have ins1:  $\bigwedge x. x \in \text{inside} (f \text{' path-image } c) \implies x \in f \text{' inside} (\text{path-image } c)$ 
proof -
  fix x

```

```

assume *:  $x \in \text{inside } (f \text{ ' path-image } c)$ 
show  $x \in f \text{ ' inside } (\text{path-image } c)$ 
  by (metis (no-types) * same-set inside-def)
qed
then have  $\text{inside } (f \text{ ' } (\text{path-image } c)) \subseteq f \text{ ' } (\text{inside}(\text{path-image } c))$ 
  by auto
have ins2:  $\bigwedge xa. xa \in \text{inside } (\text{path-image } c) \implies f \text{ ' } xa \in \text{inside } (f \text{ ' path-image } c)$ 
proof -
  fix xa
  assume *:  $xa \in \text{inside } (\text{path-image } c)$ 
  show  $f \text{ ' } xa \in \text{inside } (f \text{ ' path-image } c)$ 
    by (metis (no-types, lifting) * same-set assms(3) bij-def inj-image-mem-iff
inside-def mem-Collect-eq)
  qed
then have  $f \text{ ' } (\text{inside}(\text{path-image } c)) \subseteq \text{inside } (f \text{ ' } (\text{path-image } c))$ 
  by auto
show ?thesis
using ins1 ins2 by auto
qed

```

lemma *bij-image-intersection*:

```

assumes  $\text{path-image } c1 \cap \text{path-image } c2 = S$ 
assumes bij f
assumes  $c \in \text{path-image } (f \circ c1) \cap \text{path-image } (f \circ c2)$ 
shows  $c \in f \text{ ' } S$ 
proof -
  have  $c \in f \text{ ' path-image } c1 \cap f \text{ ' path-image } c2$ 
    using assms path-image-compose[of f c1] path-image-compose[of f c2]
    by auto
  then obtain w where c-is:  $w \in \text{path-image } c1 \wedge w \in \text{path-image } c2 \wedge c = f$ 
w
    using assms unfolding bij-def inj-def surj-def
    by auto
  then have  $w \in S$ 
    using assms by auto
  then show  $c \in f \text{ ' } S$ 
    using c-is by auto
qed

```

theorem (in *c1-on-open-R2*) *split-inside-simple-closed-curve-locale*:

```

fixes c :: real  $\implies 'a$ 
assumes c1-simple: simple-path c1 and c1-start:  $\text{pathstart } c1 = a$  and c1-end:
pathfinish c1 = b
assumes c2-simple: simple-path c2 and c2-start:  $\text{pathstart } c2 = a$  and c2-end:
pathfinish c2 = b
assumes c-simple: simple-path c and c-start:  $\text{pathstart } c = a$  and c-end: pathfin-
ish c = b
assumes a-neq-b:  $a \neq b$ 

```

and $c1c2$: $path\text{-}image\ c1 \cap path\text{-}image\ c2 = \{a,b\}$
and $c1c$: $path\text{-}image\ c1 \cap path\text{-}image\ c = \{a,b\}$
and $c2c$: $path\text{-}image\ c2 \cap path\text{-}image\ c = \{a,b\}$
and $ne\text{-}12$: $path\text{-}image\ c \cap inside(path\text{-}image\ c1 \cup path\text{-}image\ c2) \neq \{\}$
obtains $inside(path\text{-}image\ c1 \cup path\text{-}image\ c) \cap inside(path\text{-}image\ c2 \cup path\text{-}image\ c) = \{\}$
 $inside(path\text{-}image\ c1 \cup path\text{-}image\ c) \cup inside(path\text{-}image\ c2 \cup path\text{-}image\ c) \cup$
 $(path\text{-}image\ c - \{a,b\}) = inside(path\text{-}image\ c1 \cup path\text{-}image\ c2)$

proof –

let $?cc1 = (complex\text{-}of \circ c1)$
let $?cc2 = (complex\text{-}of \circ c2)$
let $?cc = (complex\text{-}of \circ c)$
have $cc1\text{-}simple$: $simple\text{-}path\ ?cc1$
using $bij\text{-}betw\text{-}imp\text{-}inj\text{-}on\ c1\text{-}simple\ complex\text{-}of\text{-}bij$
using $simple\text{-}path\text{-}linear\text{-}image\text{-}eq[OF\ complex\text{-}of\text{-}linear]$
by $blast$
have $cc1\text{-}start$: $pathstart\ ?cc1 = (complex\text{-}of\ a)$
using $c1\text{-}start$ **by** $(simp\ add: pathstart\text{-}compose)$
have $cc1\text{-}end$: $pathfinish\ ?cc1 = (complex\text{-}of\ b)$
using $c1\text{-}end$ **by** $(simp\ add: pathfinish\text{-}compose)$
have $cc2\text{-}simple$: $simple\text{-}path\ ?cc2$
using $c2\text{-}simple\ complex\text{-}of\text{-}bij\ bij\text{-}betw\text{-}imp\text{-}inj\text{-}on$
using $simple\text{-}path\text{-}linear\text{-}image\text{-}eq[OF\ complex\text{-}of\text{-}linear]$
by $blast$
have $cc2\text{-}start$: $pathstart\ ?cc2 = (complex\text{-}of\ a)$
using $c2\text{-}start$ **by** $(simp\ add: pathstart\text{-}compose)$
have $cc2\text{-}end$: $pathfinish\ ?cc2 = (complex\text{-}of\ b)$
using $c2\text{-}end$ **by** $(simp\ add: pathfinish\text{-}compose)$
have $cc\text{-}simple$: $simple\text{-}path\ ?cc$ **using** $c\text{-}simple\ complex\text{-}of\text{-}bij$
using $bij\text{-}betw\text{-}imp\text{-}inj\text{-}on$
using $simple\text{-}path\text{-}linear\text{-}image\text{-}eq[OF\ complex\text{-}of\text{-}linear]$
by $blast$
have $cc\text{-}start$: $pathstart\ ?cc = (complex\text{-}of\ a)$
using $c\text{-}start$ **by** $(simp\ add: pathstart\text{-}compose)$
have $cc\text{-}end$: $pathfinish\ ?cc = (complex\text{-}of\ b)$
using $c\text{-}end$ **by** $(simp\ add: pathfinish\text{-}compose)$
have $ca\text{-}neq\text{-}cb$: $complex\text{-}of\ a \neq complex\text{-}of\ b$
using $a\text{-}neq\text{-}b$
by $(meson\ bij\text{-}betw\text{-}imp\text{-}inj\text{-}on\ complex\text{-}of\text{-}bij\ inj\text{-}eq)$
have $image\text{-}set\text{-}eq1$: $\{complex\text{-}of\ a,\ complex\text{-}of\ b\} \subseteq path\text{-}image\ ?cc1 \cap path\text{-}image\ ?cc2$
using $c1c2\ path\text{-}image\text{-}compose[of\ complex\text{-}of\ c1]\ path\text{-}image\text{-}compose[of\ complex\text{-}of\ c2]$
by $auto$
have $image\text{-}set\text{-}eq2$: $\bigwedge c. c \in path\text{-}image\ ?cc1 \cap path\text{-}image\ ?cc2 \implies c \in \{complex\text{-}of\ a,\ complex\text{-}of\ b\}$
using $bij\text{-}image\text{-}intersection[of\ c1\ c2\ \{a,\ b\}\ complex\text{-}of]$
using $c1c2\ complex\text{-}of\text{-}bij$ **by** $auto$

```

have cc1c2: path-image ?cc1  $\cap$  path-image ?cc2 = {(complex-of a),(complex-of
b)}
  using image-set-eq1 image-set-eq2 by auto
have image-set-eq1: {complex-of a, complex-of b}  $\subseteq$  path-image ?cc1  $\cap$  path-image
?cc
  using c1c path-image-compose[of complex-of c1] path-image-compose[of com-
plex-of c]
  by auto
have image-set-eq2:  $\bigwedge c. c \in$  path-image ?cc1  $\cap$  path-image ?cc  $\implies c \in$ {complex-of
a, complex-of b}
  using bij-image-intersection[of c1 c {a, b} complex-of]
  using c1c complex-of-bij by auto
have cc1c: path-image ?cc1  $\cap$  path-image ?cc = {(complex-of a),(complex-of b)}

  using image-set-eq1 image-set-eq2 by auto
have image-set-eq1: {complex-of a, complex-of b}  $\subseteq$  path-image ?cc2  $\cap$  path-image
?cc
  using c2c path-image-compose[of complex-of c2] path-image-compose[of com-
plex-of c]
  by auto
have image-set-eq2:  $\bigwedge c. c \in$  path-image ?cc2  $\cap$  path-image ?cc  $\implies c \in$ {complex-of
a, complex-of b}
  using bij-image-intersection[of c2 c {a, b} complex-of]
  using c2c complex-of-bij by auto
have cc2c: path-image ?cc2  $\cap$  path-image ?cc = {(complex-of a),(complex-of b)}
  using image-set-eq1 image-set-eq2 by auto

let ?j = c1 +++ (reversepath c)
let ?cj = ?cc1 +++ (reversepath ?cc)
have cj-and-j: path-image ?cj = complex-of ' (path-image ?j)
  by (metis path-compose-join path-compose-reversepath path-image-compose)
have pathstart (reversepath c) = b
  using c-end
  by auto
then have j-path: path (c1 +++ (reversepath c))
  using c1-end c1-simple c-simple unfolding simple-path-def path-def
  by (metis continuous-on-joinpaths path-def path-reversepath)
then have path ?j  $\wedge$  path-image ?j = path-image c1  $\cup$  path-image c
  using <pathstart (reversepath c) = b> c1-end path-image-join path-image-reversepath
by blast
then have inside(path-image c1  $\cup$  path-image c) = inside(path-image ?j)
  by auto
have pathstart (reversepath ?cc) = complex-of b
  using cc-end
  by auto
then have cj-path: path ?cj
  using cc1-end cc1-simple cc-simple unfolding simple-path-def path-def
  by (metis continuous-on-joinpaths path-def path-reversepath)

```

then have $\text{path } ?cj \wedge \text{path-image } ?cj = \text{path-image } ?cc1 \cup \text{path-image } ?cc$
by (*metis* $\langle \text{pathstart } (\text{reversepath } (\text{complex-of } \circ c)) = \text{complex-of } b \rangle$ *cc1-end*
path-image-join path-image-reversepath)
then have $\text{ins-cj: } \text{inside}(\text{path-image } ?cc1 \cup \text{path-image } ?cc) = \text{inside } (\text{path-image } ?cj)$
by *auto*
have $\text{inside}(\text{path-image } ?cj) = \text{complex-of } ' (\text{inside}(\text{path-image } ?j))$
using *inside-bijective-linear-image*[of $?j$ *complex-of*] *j-path*
using *cj-and-j complex-of-bij complex-of-linear* **by** *presburger*
then have $i1: \text{inside}(\text{path-image } ?cc1 \cup \text{path-image } ?cc) = \text{complex-of } ' (\text{inside}(\text{path-image } c1 \cup \text{path-image } c))$ **using** *complex-of-real-of unfolding image-comp*
using *cj-and-j*
by (*simp add: ins-cj* $\langle \text{inside } (\text{path-image } c1 \cup \text{path-image } c) = \text{inside } (\text{path-image } (c1 \text{ +++ } \text{reversepath } c)) \rangle$)

let $?j2 = c2 \text{ +++ } (\text{reversepath } c)$
let $?cj2 = ?cc2 \text{ +++ } (\text{reversepath } ?cc)$
have $\text{cj2-and-j2: } \text{path-image } ?cj2 = \text{complex-of } ' (\text{path-image } ?j2)$
by (*metis* *path-compose-join path-compose-reversepath path-image-compose*)
have $\text{pathstart } (\text{reversepath } c) = b$
using *c-end by auto*
then have $\text{j2-path: } \text{path } (c2 \text{ +++ } (\text{reversepath } c))$
using *c2-end c2-simple c-simple unfolding simple-path-def path-def*
by (*metis* *continuous-on-joinpaths path-def path-reversepath*)
then have $\text{path } ?j2 \wedge \text{path-image } ?j2 = \text{path-image } c2 \cup \text{path-image } c$
using $\langle \text{pathstart } (\text{reversepath } c) = b \rangle$ *c2-end path-image-join path-image-reversepath*
by *blast*
then have $\text{inside}(\text{path-image } c2 \cup \text{path-image } c) = \text{inside}(\text{path-image } ?j2)$
by *auto*
have $\text{pathstart } (\text{reversepath } ?cc) = \text{complex-of } b$
using *cc-end by auto*
then have $\text{cj2-path: } \text{path } ?cj2$
using *cc2-end cc2-simple cc-simple unfolding simple-path-def path-def*
by (*metis* *continuous-on-joinpaths path-def path-reversepath*)
then have $\text{path } ?cj2 \wedge \text{path-image } ?cj2 = \text{path-image } ?cc2 \cup \text{path-image } ?cc$
by (*metis* $\langle \text{pathstart } (\text{reversepath } (\text{complex-of } \circ c)) = \text{complex-of } b \rangle$ *cc2-end*
path-image-join path-image-reversepath)
then have $\text{ins-cj2: } \text{inside}(\text{path-image } ?cc2 \cup \text{path-image } ?cc) = \text{inside } (\text{path-image } ?cj2)$
by *auto*
have $\text{inside}(\text{path-image } ?cj2) = \text{complex-of } ' (\text{inside}(\text{path-image } ?j2))$
using *inside-bijective-linear-image*[of $?j2$ *complex-of*] *j2-path*
using *cj2-and-j2 complex-of-bij complex-of-linear*
by *presburger*
then have $i2: \text{inside } (\text{path-image } (\text{complex-of } \circ c2) \cup \text{path-image } (\text{complex-of } \circ c))$
 $= \text{complex-of } ' \text{inside } (\text{path-image } c2 \cup \text{path-image } c)$
using *cj2-and-j2*

by (simp add: ins-cj2 <inside (path-image c2 \cup path-image c) = inside (path-image (c2 +++ reversepath c))>)

```

let ?j3 = c2 +++ (reversepath c1)
let ?cj3 = ?cc2 +++ (reversepath ?cc1)
have cj3-and-j3: path-image ?cj3 = complex-of ' (path-image ?j3)
  by (metis path-compose-join path-compose-reversepath path-image-compose)
have pathstart (reversepath c1) = b
  using c1-end by auto
then have j3-path: path (c2 +++ (reversepath c1))
  using c2-end c2-simple c1-simple unfolding simple-path-def path-def
  by (metis continuous-on-joinpaths path-def path-reversepath)
then have path-j3: path ?j3  $\wedge$  path-image ?j3 = path-image c2  $\cup$  path-image c1
  using <pathstart (reversepath c1) = b> c2-end path-image-join path-image-reversepath
by blast
then have inside(path-image c2  $\cup$  path-image c1) = inside(path-image ?j3)
  by auto
have pathstart (reversepath ?cc1) = complex-of b
  using cc1-end by auto
then have cj3-path: path ?cj3
  using cc2-end cc2-simple cc1-simple unfolding simple-path-def path-def
  by (metis continuous-on-joinpaths path-def path-reversepath)
then have path-cj3: path ?cj3  $\wedge$  path-image ?cj3 = path-image ?cc2  $\cup$  path-image
?cc1
  by (metis <pathstart (reversepath (complex-of  $\circ$  c1)) = complex-of b> cc2-end
path-image-join path-image-reversepath)
then have ins-cj3: inside(path-image ?cc2  $\cup$  path-image ?cc1) = inside (path-image
?cj3)
  by auto
have inside(path-image ?cj3) = complex-of ' (inside(path-image ?j3))
  using inside-bijective-linear-image[of ?j3 complex-of] j3-path
  using cj3-and-j3 complex-of-bij complex-of-linear
  by presburger
then have i3: inside (path-image (complex-of  $\circ$  c1)  $\cup$  path-image (complex-of  $\circ$ 
c2))
  = complex-of ' inside (path-image c1  $\cup$  path-image c2)
  by (simp add: path-cj3 path-j3 sup-commute)
obtain y where y-prop: y  $\in$  path-image c  $\cap$  inside (path-image c1  $\cup$  path-image
c2)
  using ne-12 by auto
then have y-in1: complex-of y  $\in$  path-image ?cc
  by (metis IntD1 image-eqI path-image-compose)
have y-in2: complex-of y  $\in$  complex-of ' (inside (path-image c1  $\cup$  path-image
c2))
  using y-prop by auto
then have cne-12: path-image ?cc  $\cap$  inside(path-image ?cc1  $\cup$  path-image ?cc2)
 $\neq$  {}
  using ne-12 y-in1 y-in2 i3 by force

```

obtain *for-reals*: $inside(path-image ?cc1 \cup path-image ?cc) \cap inside(path-image ?cc2 \cup path-image ?cc) = \{\}$
 $inside(path-image ?cc1 \cup path-image ?cc) \cup inside(path-image ?cc2 \cup path-image ?cc) \cup$
 $(path-image ?cc - \{complex-of a, complex-of b\}) = inside(path-image ?cc1 \cup path-image ?cc2)$
using *split-inside-simple-closed-curve*[*OF cc1-simple cc1-start cc1-end cc2-simple cc2-start*
 $cc2-end cc-simple cc-start cc-end ca-neq-cb cc1c2 cc1c cc2c cne-12$]
by *auto*
let $?rin1 = real-of ' inside(path-image ?cc1 \cup path-image ?cc)$
let $?rin2 = real-of ' inside(path-image ?cc2 \cup path-image ?cc)$

have $h1: inside(path-image c1 \cup path-image c) \cap inside(path-image c2 \cup path-image c) \neq \{\} \implies False$
proof –
assume $inside(path-image c1 \cup path-image c) \cap inside(path-image c2 \cup path-image c) \neq \{\}$
then obtain a where $a-prop: a \in inside(path-image c1 \cup path-image c) \wedge a \in inside(path-image c2 \cup path-image c)$
by *auto*
have $in1: complex-of a \in inside (path-image (complex-of \circ c1) \cup path-image (complex-of \circ c))$
using $a-prop i1$ **by** *auto*
have $in2: complex-of a \in inside (path-image (complex-of \circ c2) \cup path-image (complex-of \circ c))$
using $a-prop i2$ **by** *auto*
show $False$ **using** $in1 in2$ *for-reals(1)* **by** *auto*
qed
have $h: path-image (complex-of \circ c) - \{complex-of a, complex-of b\} = complex-of (path-image c) - complex-of \{a, b\}$
using *path-image-compose* **by** *auto*
have $complex-of ' path-image c - complex-of \{a, b\} = complex-of ' (path-image c - \{a, b\})$
proof –
have $\bigwedge x. x \in (complex-of ' path-image c - complex-of \{a, b\}) \iff x \in complex-of ' (path-image c - \{a, b\})$
using *Diff-iff bij-betw-imp-inj-on complex-of-bij image-iff inj-eq* **by** (*smt (z3)*)
then show $?thesis$ **by** *blast*
qed
then have $path-image (complex-of \circ c) - \{complex-of a, complex-of b\} = complex-of ' (path-image c - \{a, b\})$
using h **by** *simp*
then have $h2: inside(path-image c1 \cup path-image c) \cup inside(path-image c2 \cup path-image c) \cup$
 $(path-image c - \{a, b\}) = inside(path-image c1 \cup path-image c2)$
proof –
have $\bigwedge x . x \in inside(path-image c1 \cup path-image c2) \iff complex-of x \in complex-of ' inside (path-image c1 \cup path-image c2)$

using *i3* **by** (*metis bij-betw-imp-inj-on complex-of-bij image-iff inj-eq*)
then have *in-iff*: $\bigwedge x. x \in \text{inside}(\text{path-image } c1 \cup \text{path-image } c2) \longleftrightarrow \text{complex-of } x \in \text{inside}(\text{path-image}(\text{complex-of} \circ c1) \cup \text{path-image}(\text{complex-of} \circ c))$
 \cup
 $\text{inside}(\text{path-image}(\text{complex-of} \circ c2) \cup \text{path-image}(\text{complex-of} \circ c)) \cup$
 $(\text{path-image}(\text{complex-of} \circ c) - \{\text{complex-of } a, \text{complex-of } b\})$
using *for-reals(2)*
using *i3* **by** *presburger*
have $\bigwedge x. \text{complex-of } x \in \text{inside}(\text{path-image}(\text{complex-of} \circ c1) \cup \text{path-image}(\text{complex-of} \circ c)) \cup$
 $\text{inside}(\text{path-image}(\text{complex-of} \circ c2) \cup \text{path-image}(\text{complex-of} \circ c)) \cup$
 $(\text{path-image}(\text{complex-of} \circ c) - \{\text{complex-of } a, \text{complex-of } b\})$
 $\longleftrightarrow \text{complex-of } x \in \text{inside}(\text{path-image}(\text{complex-of} \circ c1) \cup \text{path-image}(\text{complex-of} \circ c))$
 $\vee \text{complex-of } x \in \text{inside}(\text{path-image}(\text{complex-of} \circ c2) \cup \text{path-image}(\text{complex-of} \circ c))$
 $\vee \text{complex-of } x \in (\text{path-image}(\text{complex-of} \circ c) - \{\text{complex-of } a, \text{complex-of } b\})$
by *blast*
then have $\bigwedge x. \text{complex-of } x \in \text{inside}(\text{path-image}(\text{complex-of} \circ c1) \cup \text{path-image}(\text{complex-of} \circ c)) \cup$
 $\text{inside}(\text{path-image}(\text{complex-of} \circ c2) \cup \text{path-image}(\text{complex-of} \circ c)) \cup$
 $(\text{path-image}(\text{complex-of} \circ c) - \{\text{complex-of } a, \text{complex-of } b\})$
 $\longleftrightarrow x \in \text{inside}(\text{path-image } c1 \cup \text{path-image } c) \cup \text{inside}(\text{path-image } c2 \cup$
 $\text{path-image } c) \cup$
 $(\text{path-image } c - \{a, b\})$
using *i1 i2 i3 Un-iff* $\langle \text{path-image}(\text{complex-of} \circ c) - \{\text{complex-of } a, \text{complex-of } b\} = \text{complex-of} \circ (\text{path-image } c - \{a, b\}) \rangle$ *bij-betw-imp-inj-on complex-of-bij image-iff inj-def*
by (*smt (verit, best)*)
then have $\bigwedge x. x \in \text{inside}(\text{path-image } c1 \cup \text{path-image } c2) \longleftrightarrow x \in (\text{inside}(\text{path-image } c1 \cup \text{path-image } c) \cup \text{inside}(\text{path-image } c2 \cup \text{path-image } c) \cup$
 $(\text{path-image } c - \{a, b\}))$
using *in-iff* **by** *meson*
then show *?thesis* **by** *auto*
qed
show *?thesis* **using** *that h1 h2* **by** *auto*
qed

lemma *split-inside-simple-closed-curve-real2*:

fixes $c :: \text{real} \Rightarrow \text{real}^2$
assumes *c1-simple*: *simple-path* $c1$ **and** *c1-start*: *pathstart* $c1 = a$ **and** *c1-end*: *pathfinish* $c1 = b$
assumes *c2-simple*: *simple-path* $c2$ **and** *c2-start*: *pathstart* $c2 = a$ **and** *c2-end*: *pathfinish* $c2 = b$
assumes *c-simple*: *simple-path* c **and** *c-start*: *pathstart* $c = a$ **and** *c-end*: *pathfinish* $c = b$
assumes *a-neq-b*: $a \neq b$
and *c1c2*: $\text{path-image } c1 \cap \text{path-image } c2 = \{a, b\}$

```

    and c1c: path-image c1 ∩ path-image c = {a,b}
    and c2c: path-image c2 ∩ path-image c = {a,b}
    and ne-12: path-image c ∩ inside(path-image c1 ∪ path-image c2) ≠ {}
obtains inside(path-image c1 ∪ path-image c) ∩ inside(path-image c2 ∪ path-image
c) = {}
    inside(path-image c1 ∪ path-image c) ∪ inside(path-image c2 ∪ path-image
c) ∪
    (path-image c - {a,b}) = inside(path-image c1 ∪ path-image c2)
proof -
  have good-type: c1-on-open-R2-axioms TYPE((real, 2) vec)
  unfolding c1-on-open-R2-axioms-def by auto
  then show ?thesis
  using c1-on-open-R2.split-inside-simple-closed-curve-locale[of - - - c1 a b c2 c]
  assms
  unfolding c1-on-open-R2-def c1-on-open-euclidean-def c1-on-open-def
  using good-type that by blast
qed

end
theory Polygon-Lemmas
imports
  Polygon-Jordan-Curve
  HOL-Library.Sublist
  HOL.Set-Interval
  HOL.Fun

```

begin

5 Properties of make polygonal path, pathstart and pathfinish of a polygon

lemma *make-polygonal-path-induct*[case-names *Empty Single Two Multiple*]:

```

fixes ell :: (real^2) list
assumes empty:  $\bigwedge ell. ell = [] \implies P ell$ 
  and single:  $\bigwedge ell. \llbracket length\ ell = 1 \rrbracket \implies P ell$ 
  and two:  $\bigwedge ell. \llbracket length\ ell = 2 \rrbracket \implies P ell$ 
  and multiple:  $\bigwedge ell. \llbracket length\ ell > 2; P ((ell!0), (ell!1)); P ((ell!1)\#(drop\ 2\ ell)) \rrbracket \implies P ell$ 
shows  $P ell$ 
apply(induct ell rule: make-polygonal-path.induct)
using empty single two multiple by auto

```

lemma *make-polygonal-path-gives-path*:

```

fixes v :: (real^2) list
shows path (make-polygonal-path v)
proof(induction length v arbitrary: v)

```

```

case 0
thus path (make-polygonal-path v)
  by auto
next
case (Suc x)
show ?case
  by (smt (verit, best) Suc.hyps(1) Suc.hyps(2) Suc-length-conv list.distinct(1)
list.inject make-polygonal-path.elims path-join-imp path-linepath pathfinish-linepath
pathstart-join pathstart-linepath)
qed

```

```

corollary polygonal-path-is-path:
  fixes g :: R-to-R2
  assumes polygonal-path g
  shows path g
  using assms polygonal-path-def make-polygonal-path-gives-path by auto

```

```

lemma polygon-to-polygonal-path:
  fixes p :: R-to-R2
  assumes polygon p
  obtains ell where p = make-polygonal-path ell
  using assms unfolding polygon-def polygonal-path-def
  by auto

```

```

lemma polygon-pathstart:
  fixes g :: R-to-R2
  assumes l ≠ []
  assumes g = make-polygonal-path l
  shows pathstart g = l!0
  using assms make-polygonal-path.simps
  by (smt (verit) list.discI list.expand make-polygonal-path.elims nth-Cons-0 path-
start-join pathstart-linepath)

```

```

lemma polygon-pathfinish:
  fixes g :: R-to-R2
  assumes l ≠ []
  assumes g = make-polygonal-path l
  shows pathfinish g = l!(length l - 1)
  using assms
proof (induct length l arbitrary: g l)
  case 0
  then show ?case by auto
next
case (Suc x)
  {assume *: length l = 1
  then obtain a where l-is: l = [a]
  by (metis Suc.prem(1) Suc-neq-Zero diff-Suc-1 diff-self-eq-0 length-Cons
remdups-adj.cases)}

```

```

then have pathfinish g = a
  using Suc make-polygonal-path.simps
  by (simp add: pathfinish-def)
then have pathfinish g = l!(length l - 1)
  using Suc l-is
  by auto
} moreover {assume *: length l = 2
  then obtain a b where l-is: l = [a, b]
    by (metis (no-types, opaque-lifting) One-nat-def Suc-eq-plus1 list.size(3)
list.size(4) min-list.cases nat.simps(1) nat.simps(3) numeral-2-eq-2)
  then have g-is: g = linepath a b
    using Suc by auto
  have pf: pathfinish g = b using g-is by auto
  then have pathfinish g = l!(length l - 1)
    using Suc * l-is
    by auto
}
moreover {assume *: length l > 2
  then obtain a b c where l-is: l = a # b # c
    by (metis Suc.prem(1) Zero-neq-Suc length-Cons less-Suc0 list.size(3)
numeral-2-eq-2 remdups-adj.cases)
  then have g-is: g = (linepath a b) +++ make-polygonal-path (b # c)
    using Suc l-is
  proof -
    have c ≠ []
      using * l-is by auto
    then show ?thesis
      by (metis (full-types) Suc(4) l-is list.exhaust make-polygonal-path.simps(4))
  qed
  then have pf: pathfinish g = pathfinish (make-polygonal-path (b # c))
    by auto
  have len-x: length (b # c) = x
    using l-is Suc by auto
  then have pathfinish (make-polygonal-path (b # c)) = (b # c)!(length l - 2)
    using Suc.hyps l-is
    by simp
  then have pathfinish g = l!(length l - 1)
    using l-is pf
    by auto
}
ultimately show ?case
  using Suc
  by (metis One-nat-def less-Suc-eq-0-disj less-antisym numeral-2-eq-2)
qed

```

lemma *make-polygonal-path-image-property:*

assumes $\text{length } vts \geq 2$

assumes $p\text{-is-path: } x \in \text{path-image } (\text{make-polygonal-path } vts)$

shows $\exists k < \text{length } vts - 1. x \in \text{path-image } (\text{linepath } (vts ! k) (vts ! (k + 1)))$

```

using assms
proof (induct vts)
  case Nil
  then show ?case by auto
next
  case (Cons a vts)
  then have len-gteq: length vts ≥ 1
    by simp
  {assume *: length vts = 1
    then obtain b where vts-is: vts = [b]
    by (metis One-nat-def <1 ≤ length vts> drop-eq-Nil id-take-nth-drop less-numeral-extra(1)
self-append-conv2 take-eq-Nil2)
    then have x ∈ path-image (make-polygonal-path [a, b])
      using Cons by auto
    then have x ∈ path-image (linepath a b)
      by auto
    then have x ∈ path-image (linepath ((a#vts) ! 0) ((a#vts) ! 1))
      using Cons vts-is
      by force
    then have  $\exists k < \text{length } (a \# vts) - 1. x \in \text{path-image } (\text{linepath } ((a \# vts) ! k) ((a \# vts) ! (k + 1)))$ 
      using *
      by simp
  } moreover {assume *: length vts > 1
    then obtain b vts' where vts-is: vts = b # vts'
    by (metis One-nat-def le-zero-eq len-gteq list.exhaust list.size(3) n-not-Suc-n)
    then have x ∈ path-image ((linepath a b) +++ make-polygonal-path (b # vts'))
      using Cons
    by (metis (no-types, lifting) * One-nat-def length-Cons list.exhaust list.size(3)
make-polygonal-path.simps(4) nat-less-le)
    then have eo: x ∈ path-image ((linepath a b)) ∨ x ∈ path-image (make-polygonal-path (b # vts'))
      using not-in-path-image-join by blast
    {assume **: x ∈ path-image ((linepath a b))
      then have  $\exists k < \text{length } (a \# vts) - 1. x \in \text{path-image } (\text{linepath } ((a \# vts) ! k) ((a \# vts) ! (k + 1)))$ 
        using vts-is
        by auto
    } moreover {assume **: x ∈ path-image (make-polygonal-path (b # vts'))
      then have  $\exists k < \text{length } vts - 1. x \in \text{path-image } (\text{linepath } (vts ! k) (vts ! (k + 1)))$ 
        using Cons.hyps(1) *
        by (simp add: Suc-leI vts-is)
      then have  $\exists k < \text{length } (a \# vts) - 1. x \in \text{path-image } (\text{linepath } ((a \# vts) ! k) ((a \# vts) ! (k + 1)))$ 
    }

    using add commute add-diff-cancel-left' length-Cons less-diff-conv nth-Cons-Suc plus-1-eq-Suc by auto
  }

```

```

    ultimately have  $\exists k < \text{length } (a \# \text{vts}) - 1. x \in \text{path-image } (\text{linepath } ((a \# \text{vts}) ! k) ((a \# \text{vts}) ! (k + 1)))$ 
      using eo by auto
  }
  ultimately show ?case
    using len-gteq
    by fastforce
qed

lemma linepaths-subset-make-polygonal-path-image:
  assumes length vts  $\geq 2$ 
  assumes k  $< \text{length } \text{vts} - 1$ 
  shows path-image (linepath (vts ! k) (vts ! (k + 1)))  $\subseteq$  path-image (make-polygonal-path vts)
  using assms
proof (induct vts arbitrary: k)
  case Nil
  then show ?case by auto
next
  case (Cons a vts)
  { assume *: length vts = 1
    then have k-is: k = 0
      using Cons.prem2 by auto
    obtain b where vts-is: vts = [b]
      using *
    by (metis One-nat-def drop-eq-Nil id-take-nth-drop le-numeral-extra(4) self-append-conv2 take-eq-Nil2 zero-less-one)
    then have path-image (make-polygonal-path (a # vts)) = path-image (linepath a b)
      by auto
    then have path-image (linepath ((a # vts) ! k) ((a # vts) ! (k + 1)))
       $\subseteq$  path-image (make-polygonal-path (a # vts))
      using k-is vts-is
      by simp
  } moreover
  { assume *: length vts  $> 1$ 
    then obtain b c vts' where vts-is: vts = b # c # vts'
      by (metis diff-0-eq-0 diff-Suc-1 diff-is-0-eq leD length-Cons list.exhaust list.size(3))
    { assume **: k = 0
      then have same-path-image: path-image (linepath ((a # vts) ! k) ((a # vts) ! (k + 1))) = path-image (linepath a b)
        using vts-is
        by auto
      have path-image (linepath a b)  $\subseteq$  path-image (make-polygonal-path (a # b # c # vts'))
        using vts-is make-polygonal-path.simps path-image-join
        by (metis (no-types, lifting) Un-iff list.discI nth-Cons-0 pathfinish-linepath polygon-pathstart subsetI)
      then have path-image (linepath ((a # vts) ! k) ((a # vts) ! (k + 1)))  $\subseteq$ 

```

```

path-image (make-polygonal-path (a # vts))
  using vts-is same-path-image
  by presburger
} moreover {assume **: k > 0
  then have k-minus-lt: k-1 < length vts - 1
    using Cons
    by auto
  then have path-image-is: path-image (linepath ((a # vts) ! k) ((a # vts) ! (k
+ 1))) = path-image (linepath (vts ! (k-1)) (vts ! k))
    using **
    by auto
  then have path-im-subset1: path-image (linepath (vts ! (k-1)) (vts ! k)) ⊆
path-image (make-polygonal-path vts)
    using k-minus-lt Cons.hyps(1)[of k-1] * ** Suc-leI Suc-pred add.right-neutral
add-Suc-right nat-1-add-1 plus-1-eq-Suc
    by auto
  have path-im-subset2: path-image (make-polygonal-path vts) ⊆ path-image
(make-polygonal-path (a # vts))
    using vts-is make-polygonal-path.simps(4)
    by (metis dual-order.refl list.distinct(1) nth-Cons-0 path-image-join pathfin-
ish-linepath polygon-pathstart sup.coboundedI2)
  then have path-image (linepath ((a # vts) ! k) ((a # vts) ! (k + 1))) ⊆
path-image (make-polygonal-path (a # vts))
    using path-image-is path-im-subset1 path-im-subset2
    by blast
}
ultimately have path-image (linepath ((a # vts) ! k) ((a # vts) ! (k + 1)))
⊆ path-image (make-polygonal-path (a # vts))
  by blast
}
ultimately show ?case
  by (metis Cons.prem(1) Suc-1 leD length-Cons linorder-neqE-nat nat-add-left-cancel-less
plus-1-eq-Suc)
qed

```

lemma vertices-on-path-image: shows $set\ vts \subseteq path\ image\ (make\ polygonal\ path\ vts)$

proof (induct vts rule:make-polygonal-path.induct)

case 1

then show ?case by auto

next

case (2 a)

then show ?case by auto

next

case (3 a b)

then show ?case by auto

next

case (4 a b v va)

then have a-in-image: $a \in path\ image\ (make\ polygonal\ path\ (a\ \#\ b\ \#\ v\ \#\ va))$

```

using make-polygonal-path.simps
by (metis list.distinct(1) nth-Cons-0 pathstart-in-path-image polygon-pathstart)

have path-image-union:
  path-image (make-polygonal-path (a # b # v # va))
  = path-image (linepath a b)  $\cup$  path-image (make-polygonal-path (b # v # va))
by (metis make-polygonal-path.simps(4) linepath-1' list.discI nth-Cons-0 path-image-join
pathfinish-def polygon-pathstart)
have set (a # b # v # va) = {a}  $\cup$  set (b # v # va)
by auto
then show ?case using a-in-image 4 make-polygonal-path.simps
path-image-union by auto
qed

lemma path-image-cons-union:
  assumes p = make-polygonal-path vts
  assumes p' = make-polygonal-path vts'
  assumes vts'  $\neq$  []
  assumes vts = a # vts'  $\wedge$  b = vts'!0
  shows path-image p = path-image (linepath a b)  $\cup$  path-image p'
proof –
  have pathfinish (linepath a b) = pathstart p' using assms polygon-pathstart by
auto
  moreover have length vts = 2  $\implies$  ?thesis
  by (smt (verit) Cons-nth-drop-Suc One-nat-def Suc-1 assms(1) assms(2) assms(3)
assms(4) closed-segment-idem diff-Suc-1 drop0 drop-eq-Nil insert-subset le-iff-sup
le-numeral-extra(4) length-Cons length-greater-0-conv list.discI list.inject list.set(1)
list.set(2) make-polygonal-path.elims path-image-linepath sup-commute vertices-on-path-image)
  moreover have length vts > 2  $\implies$  ?thesis
  by (metis (no-types, lifting) Cons-nth-drop-Suc One-nat-def Suc-1 assms(1)
assms(2) assms(3) assms(4) calculation(1) drop0 drop-Suc-Cons length-greater-0-conv
make-polygonal-path.simps(4) path-image-join)
  moreover have length vts  $\geq$  2 using assms by (simp add: Suc-le-eq)
  ultimately show ?thesis by linarith
qed

lemma polygonal-path-image-linepath-union:
  assumes p = make-polygonal-path vts
  assumes n = length vts
  assumes n  $\geq$  2
  shows path-image p = ( $\bigcup$  {path-image (linepath (vts!i) (vts!(i+1))) | i. i  $\leq$  n
  – 2})
  using assms
proof(induct n arbitrary: vts p)
  case 0
  then show ?case by linarith
next
  case (Suc n)
  { assume *: Suc n = 2

```

then obtain $a\ b$ **where** $ab: vts = [a, b]$
by (*metis Suc.premis(2-3) Cons-nth-drop-Suc One-nat-def Suc-1 drop0 drop-eq-Nil lessI pos2*)
then have $path\ image\ p = path\ image\ (linepath\ a\ b)$
using *make-polygonal-path.simps Suc.premis by presburger*
moreover have $\dots = (\bigcup \{path\ image\ (linepath\ (vts!\ i)\ (vts!\ (i+1))) \mid i.\ i \leq Suc\ n - 2\})$
using *ab Suc.premis*
by (*smt (verit, ccfv-threshold) Suc-eq-plus1 Sup-least Sup-upper * diff-is-0-eq diff-zero dual-order.refl mem-Collect-eq nth-Cons-0 nth-Cons-Suc subset-antisym*)
ultimately have *?case by presburger*
} moreover
{ assume $*$: $Suc\ n > 2$
then obtain $a\ b\ vts'$ **where** $vts': vts = a \# vts' \wedge b = vts'!0 \wedge vts' = tl\ vts$
by (*metis Suc.premis(2) list.collapse list.size(3) nat.distinct(1)*)

let $?p' = make\ polygonal\ path\ vts'$
let $?P' = path\ image\ ?p'$
let $?P = path\ image\ p$
let $?P\ union = (\bigcup \{path\ image\ (linepath\ (vts!\ i)\ (vts!\ (i+1))) \mid i.\ i \leq n - 1\})$

have $vts'\text{-len}: length\ vts' = n$ **using** $vts'\ Suc.premis$ **by** *fastforce*
then have $?P' = (\bigcup \{path\ image\ (linepath\ (vts!\ i)\ (vts'\ (i+1))) \mid i.\ i \leq n - 2\})$
using *Suc.premis Suc.hyps * by force*
moreover have $\forall i \leq n-2.\ vts'\ i = vts'\ (i+1) \wedge vts'\ (i+1) = vts'\ (i+2)$ **using** vts' **by** *force*
ultimately have $?P' = (\bigcup \{path\ image\ (linepath\ (vts'\ (i+1))\ (vts'\ (i+2))) \mid i.\ i \leq n - 2\})$
by *fastforce*
moreover have $\dots = (\bigcup \{path\ image\ (linepath\ (vts!\ i)\ (vts'\ (i+1))) \mid i.\ 1 \leq i \wedge i \leq n - 1\})$
(is $\dots = ?P'\text{-union}$ **)**
proof-
have $\bigwedge x\ i.\ x \in \{vts!\ i \text{---} vts!\ Suc\ (Suc\ i)\}$
 $\implies i \leq n - 2$
 $\implies \exists xa.\ (\exists i.\ xa = \{vts!\ i \text{---} vts!\ Suc\ i\} \wedge Suc\ 0 \leq i \wedge i \leq n - Suc\ 0)$
 $\wedge x \in xa$
by (*metis * One-nat-def Suc-diff-Suc Suc-le-mono add-2-eq-Suc' bot-nat-0.extremum diff-Suc-Suc le-add-diff-inverse plus-1-eq-Suc*)
moreover have $\bigwedge x\ i.\ x \in \{vts!\ i \text{---} vts!\ Suc\ i\}$
 $\implies Suc\ 0 \leq i$
 $\implies i \leq n - Suc\ 0$
 $\implies \exists xa.\ (\exists i.\ xa = \{vts!\ Suc\ i \text{---} vts!\ Suc\ (Suc\ i)\} \wedge i \leq n - 2) \wedge x \in xa$
by (*metis * Suc-diff-Suc gr0-implies-Suc linorder-not-le not-less-eq-eq numeral-2-eq-2*)
ultimately show *?thesis by auto*
qed

moreover have $\text{path-image } (\text{linepath } a \ b) \cup ?P' \text{-union} = ?P \text{-union}$
proof –
have $\bigwedge x. x \in \{a \ -- \ b\} \implies \exists xa. (\exists i. xa = \{\text{vts} \ ! \ i \ -- \ \text{vts} \ ! \ \text{Suc } i\} \wedge i \leq n - \text{Suc } 0) \wedge x \in xa$
using vts' **by** fastforce
moreover have $\bigwedge x i. x \in \{\text{vts} \ ! \ i \ -- \ \text{vts} \ ! \ \text{Suc } i\}$
 $\implies \forall xa. (\forall i \geq \text{Suc } 0. xa = \{\text{vts} \ ! \ i \ -- \ \text{vts} \ ! \ \text{Suc } i\} \longrightarrow \neg i \leq n - \text{Suc } 0)$
 $\vee x \notin xa$
 $\implies i \leq n - \text{Suc } 0$
 $\implies x \in \{a \ -- \ b\}$
by ($\text{metis } \text{Suc-le-eq } \text{bot-nat-0.not-eq-extremum } \text{nth-Cons-0 } \text{nth-Cons-Suc } \text{vts}'$)
ultimately show $?thesis$ **by** auto
qed
moreover have $?P = (\text{path-image } (\text{linepath } a \ b)) \cup ?P'$
using $\text{Suc.prem } \text{vts}'$ $\text{path-image-cons-union}$
by ($\text{metis } \text{One-nat-def } \text{Suc-1 } \text{vts}'\text{-len } \text{bot-nat-0.extremum } \text{list.size}(3) \text{not-less-eq-eq}$)
ultimately have $?case$ **by** force
}
ultimately show $?case$ **using** Suc.prem **by** linarith
qed

6 Loop Free Properties

lemma $\text{constant-linepath-is-not-loop-free}$:

shows $\neg(\text{loop-free } ((\text{linepath } a \ a)::\text{real} \Rightarrow \text{real}^2))$

proof –

have $\text{all-zero1}: \bigwedge x y::\text{real}. (1 - x) *_{\mathbb{R}} (a::\text{real}^2) + x *_{\mathbb{R}} a = a$

by auto

have $\text{all-zero2}: \bigwedge x y::\text{real}. (1 - y) *_{\mathbb{R}} (a::\text{real}^2) + y *_{\mathbb{R}} a = a$

by auto

then have $\exists x::\text{real} \in \{0..1\}. \exists y::\text{real} \in \{0..1\}. x \neq y \wedge (x = 0 \longrightarrow y \neq 1) \wedge (x = 1 \longrightarrow y \neq 0)$

by ($\text{metis } \text{atLeastAtMost-iff } \text{field-lbound-gt-zero } \text{less-eq-real-def } \text{linorder-not-less-zero-less-one}$)

then show $?thesis$

unfolding $\text{loop-free-def } \text{linepath-def}$

using $\text{all-zero1 } \text{all-zero2}$ **by** auto

qed

lemma $\text{doubling-back-is-not-loop-free}$:

assumes $a \neq b$

shows $\neg(\text{loop-free } ((\text{make-polygonal-path } [a, b, a])::\text{real} \Rightarrow \text{real}^2))$

proof –

let $?p1 = (1/4::\text{real})$

let $?p2 = (3/4::\text{real})$

have $\text{same-point}: ((\text{linepath } a \ b) \ +++ (\text{linepath } b \ a)) (1/4::\text{real}) = ((\text{linepath } a \ b) \ +++ (\text{linepath } b \ a)) (3/4::\text{real})$

unfolding $\text{linepath-def } \text{joinpaths-def}$ **by** auto

have $?p1 \in \{0..1\} \wedge ?p2 \in \{0..1\} \wedge ?p1 \neq ?p2 \wedge (?p1 = 0 \longrightarrow ?p2 \neq 1) \wedge$
 $(?p1 = 1 \longrightarrow ?p2 \neq 0)$
by *auto*
then have $\exists x \in \{0..1\}. \exists y \in \{0..1\}.$
 $(\text{linepath } a \ b \ +++ \ \text{linepath } b \ a) \ x = (\text{linepath } a \ b \ +++ \ \text{linepath } b \ a) \ y$
 $\wedge x \neq y \wedge (x = 0 \longrightarrow y \neq 1) \wedge (x = 1 \longrightarrow y \neq 0)$
using *same-point by blast*
then have $\neg(\text{loop-free } ((\text{linepath } a \ b) \ +++ \ (\text{linepath } b \ a)))$
unfolding *loop-free-def by auto*
then show *?thesis using make-polygonal-path.simps*
by *auto*
qed

lemma *not-loop-free-first-component:*

assumes $\neg(\text{loop-free } p1)$
shows $\neg(\text{loop-free } (p1 \ +++ \ p2))$

proof –

obtain $x \ y$ **where** *xy-prop*: $0 \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1 \ x \neq y$
 $(x = 0 \longrightarrow y \neq 1) \ (x = 1 \longrightarrow y \neq 0)$
 $p1 \ x = p1 \ y$
using *assms unfolding loop-free-def*
by *auto*

then have *xy-prop2*: $0 \leq x/2 \ x/2 \leq 1/2 \ 0 \leq y/2 \ y/2 \leq 1/2 \ x/2 \neq y/2$
by *auto*

then have $(p1 \ +++ \ p2) \ (x/2) = (p1 \ +++ \ p2) \ (y/2)$
unfolding *joinpaths-def using xy-prop(8)*
by *auto*

then have *props*: $(p1 \ +++ \ p2) \ (x/2) = (p1 \ +++ \ p2) \ (y/2) \wedge$
 $(x/2) \neq (y/2) \wedge ((x/2) = 0 \longrightarrow (y/2) \neq 1) \wedge ((x/2) = 1 \longrightarrow (y/2) \neq$
 $0)$

using *xy-prop2 by auto*

have $x/2 \in \{0..1\} \wedge y/2 \in \{0..1\}$

using *xy-prop2 by auto*

then have $\exists x \in \{0..1\}.$

$\exists y \in \{0..1\}.$

$(p1 \ +++ \ p2) \ x = (p1 \ +++ \ p2) \ y \wedge$

$x \neq y \wedge (x = 0 \longrightarrow y \neq 1) \wedge (x = 1 \longrightarrow y \neq 0)$

using *props*

by *blast*

then show *?thesis*

unfolding *loop-free-def by auto*

qed

lemma *not-loop-free-second-component:*

assumes *pathfinish-pathstart*: $\text{pathfinish } p1 = \text{pathstart } p2$

assumes $\neg(\text{loop-free } p2)$

shows $\neg(\text{loop-free } (p1 \ +++ \ p2))$

proof –

obtain $x \ y$ **where** *xy-prop*: $0 \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1 \ x \neq y$

```

    (x = 0 → y ≠ 1) (x = 1 → y ≠ 0)
  p2 x = p2 y
  using assms unfolding loop-free-def
  by auto
  then have xy-prop2: (x + 1)/2 ≥ 1/2 (x + 1)/2 ≤ 1 (y + 1)/2 ≥ 1/2 (y +
1)/2 ≤ 1
  (x + 1)/2 ≠ (y + 1)/2
  by auto
  have x-same: 2*((x + 1)/2) - 1 = x
  by (metis add.right-neutral add-diff-eq cancel-comm-monoid-add-class.diff-cancel
class-dense-linordered-field.between-same mult-1 mult-2 times-divide-eq-left times-divide-eq-right)
  have y-same: 2*((y + 1)/2) - 1 = y
  by (metis add.right-neutral add-diff-eq cancel-comm-monoid-add-class.diff-cancel
class-dense-linordered-field.between-same mult-1 mult-2 times-divide-eq-left times-divide-eq-right)
  have p2 (2*((x + 1)/2) - 1) = p2 (2*((y + 1)/2) - 1)
  using xy-prop(8) x-same y-same
  by auto
  have relate-start-finish: p1 1 = p2 0
  using pathfinish-pathstart
  unfolding pathfinish-def pathstart-def
  by auto
  then have xh1: (x + 1)/2 = 1/2 ⇒ (p1 +++ p2) ((x + 1)/2) = p2 x
  unfolding joinpaths-def
  by auto
  have xh2: (x + 1)/2 > 1/2 ⇒ (p1 +++ p2) ((x + 1)/2) = p2 x
  using xy-prop2 unfolding joinpaths-def
  using x-same by force
  then have xh: (p1 +++ p2) ((x + 1)/2) = p2 x
  using xh1 xh2 xy-prop2
  by linarith
  have yh1: (y + 1)/2 = 1/2 ⇒ (p1 +++ p2) ((y + 1)/2) = p2 y
  using relate-start-finish unfolding joinpaths-def
  by auto
  have yh2: (y + 1)/2 > 1/2 ⇒ (p1 +++ p2) ((y + 1)/2) = p2 y
  using xy-prop2 unfolding joinpaths-def
  using y-same by force
  then have yh: (p1 +++ p2) ((y + 1)/2) = p2 y
  using yh1 yh2 xy-prop2
  by linarith
  then have same-eval: (p1+++p2) ((x + 1)/2) = (p1+++p2) ((y + 1)/2)
  using xh yh xy-prop(8)
  by presburger
  have inset1: (x + 1)/2 ∈ {0..1}
  using xy-prop2
  by simp
  have inset2: (y + 1)/2 ∈ {0..1}
  using xy-prop2
  by simp
  have ∃ x∈{0..1}.

```

```

     $\exists y \in \{0..1\}$ .
     $(p1 \text{ +++ } p2) \ x = (p1 \text{ +++ } p2) \ y \wedge$ 
     $x \neq y \wedge (x = 0 \longrightarrow y \neq 1) \wedge (x = 1 \longrightarrow y \neq 0)$ 
    using xy-prop2 same-eval inset1 inset2
    by fastforce
    then show ?thesis
    unfolding loop-free-def by auto
qed

```

```

lemma loop-free-subpath:
  assumes path p
  assumes u-and-v:  $u \in \{0..1\} \ v \in \{0..1\} \ u < v$ 
  assumes  $\neg (\text{loop-free } (\text{subpath } u \ v \ p))$ 
  shows  $\neg (\text{loop-free } p)$ 
proof -
  have path (subpath u v p)
  using path-subpath assms by auto
  then show ?thesis using simple-path-subpath assms
  unfolding simple-path-def
  by blast
qed

```

```

lemma loop-free-associative:
  assumes path p
  assumes path q
  assumes path r
  assumes pathfinish p = pathstart q
  assumes pathfinish q = pathstart r
  shows  $\neg (\text{loop-free } ((p \text{ +++ } q) \text{ +++ } r)) \longleftrightarrow \neg (\text{loop-free } (p \text{ +++ } (q \text{ +++ } r)))$ 
  by (metis (mono-tags, lifting) assms(1) assms(2) assms(3) assms(4) assms(5))
path-join-imp pathfinish-join pathstart-join simple-path-assoc simple-path-def

```

```

lemma polygon-at-least-3-vertices:
  assumes polygon p and
     $p = \text{make-polygonal-path } vts$ 
  shows  $\text{card } (\text{set } vts) \geq 3$ 
  using assms
proof (induct vts rule: make-polygonal-path.induct)
  case 1
  then show ?case unfolding polygon-def
  using constant-linepath-is-not-loop-free make-polygonal-path.simps(1)
  by (metis simple-path-def)
next
  case (2 a)
  then show ?case unfolding polygon-def
  using constant-linepath-is-not-loop-free make-polygonal-path.simps(2)
  by (metis simple-path-def)
next
  case (3 a b)

```

```

{ assume *: a = b
  then have False using 3 unfolding polygon-def
    using constant-linepath-is-not-loop-free make-polygonal-path.simps(3)
    by (metis simple-path-def)
} moreover {assume *: a ≠ b
  then have False using 3 unfolding polygon-def closed-path-def
    pathstart-def pathfinish-def using make-polygonal-path.simps(3)
    by (simp add: linepath-0' linepath-1')
}
}
ultimately show ?case
  by auto
next
case (4 a b v va)
have finset: finite (set (a # b # v # va))
  by blast
have subset: {a, b, v} ⊆ set (a # b # v # va)
  by auto
have neq1: a ≠ b
  using constant-linepath-is-not-loop-free not-loop-free-first-component
  by (metis 4.prem(2) make-polygonal-path.simps(4) polygon-def assms(1) simple-path-def)
have loop-free-2: loop-free (make-polygonal-path (b # v # va))
  using 4 not-loop-free-second-component
  by (metis make-polygonal-path.simps(4) polygon-def list.distinct(1) nth-Cons-0 pathfinish-linepath polygon-pathstart simple-path-def)
have contra: b = v ⇒ ¬(loop-free (make-polygonal-path (b # v # va)))
  using constant-linepath-is-not-loop-free[of b] make-polygonal-path.simps not-loop-free-first-component
  by (metis neq-Nil-conv)
then have neq2: b ≠ v
  using loop-free-2 contra
  by auto

have ¬ loop-free ((linepath a b) +++ (linepath b a))
  using doubling-back-is-not-loop-free[of a b] neq1
  by auto
have make-path-is: make-polygonal-path (a # b # a # va) = (linepath a b) +++ ((linepath b a) +++ (make-polygonal-path (a#va)))
  using make-polygonal-path.simps
  by (metis (no-types, opaque-lifting) 4.prem(1) 4.prem(2) closed-path-def polygon-def <¬ loop-free (linepath a b +++ linepath b a)> linepath-1' min-list.cases nth-Cons-0 pathfinish-def pathfinish-join polygon-pathstart simple-path-def)
have ¬ loop-free (((linepath a b) +++ (linepath b a)) +++ (make-polygonal-path (a#va)))
  using make-polygonal-path.simps not-loop-free-first-component
  using <¬ loop-free (linepath a b +++ linepath b a)>
  by auto
then have ¬ loop-free (make-polygonal-path (a # b # a # va))
  using loop-free-associative

```

```

    by (metis make-polygonal-path-gives-path list.discI make-path-is nth-Cons-0
path-linepath pathfinish-linepath pathstart-linepath polygon-pathstart)
  then have neq3:  $v \neq a$ 
    using 4
    using polygon-def simple-path-def by blast
  have card-3:  $\text{card } \{a, b, v\} = 3$ 
    using neq1 neq2 neq3
    by auto
  then show ?case
    using subset finset
    by (metis card-mono)
qed

```

lemma *polygon-vertices-length-at-least-4*:

```

  assumes polygon p and
     $p = \text{make-polygonal-path } vts$ 
  shows  $\text{length } vts \geq 4$ 
proof -
  have card-set:  $\text{card } (\text{set } vts) \geq 3$ 
    using polygon-at-least-3-vertices assms
    by blast
  have len-gt3:  $\text{length } vts \geq 3$ 
    using card-length local.card-set order-trans by blast
  then have non-empty:  $vts \neq []$ 
    using card-set
    by auto
  have eq:  $p\ 0 = p\ 1$ 
    using assms unfolding polygon-def closed-path-def pathstart-def pathfinish-def
  by auto
  have p0:  $p\ 0 = vts\ !\ 0$ 
    using polygon-pathstart[OF non-empty] using assms unfolding pathstart-def
    by auto
  have p1:  $p\ 1 = vts\ !\ (\text{length } vts - 1)$ 
    using polygon-pathfinish[OF non-empty] using assms unfolding pathfinish-def
    by auto
  have vts ! 0 = vts ! (length vts - 1)
    using assms unfolding polygon-def
    using p0 p1 eq by auto
  then have set vts = set (drop 1 vts)
    using len-gt3
    by (smt (verit, best) Cons-nth-drop-Suc Suc-eq-plus1 Suc-le-eq add commute
add-0 add-leD2 drop0 dual-order.refl insert-subset last.simps last-conv-nth last-in-set
list.distinct(1) list.set(2) numeral-3-eq-3 order-antisym-conv)
  then have  $\text{length } (\text{drop } 1\ vts) \geq 3$ 
    using card-set
    by (metis dual-order.trans length-remdups-card-conv length-remdups-leq)
  then show ?thesis
using card-set
by (metis One-nat-def Suc-1 Suc-eq-plus1 Suc-pred add-Suc-right length-drop

```

length-greater-0-conv non-empty not-less-eq-eq numeral-3-eq-3 numeral-Bit0)
qed

lemma *linepath-loop-free*:

assumes $a \neq b$
shows *loop-free* (*linepath* a b)
unfolding *loop-free-def linepath-def*
by (*smt* ($z3$) *add.assoc add.commute add-scaleR-degen assms diff-add-cancel scaleR-left-diff-distrib*)

7 Explicit Linepath Characterization of Polygonal Paths

lemma *triangle-linepath-images*:

fixes $x :: \text{real}$
assumes $vts = [a, b, c]$
assumes $p = \text{make-polygonal-path } vts$
shows $x \in \{0..1/2\} \implies p \ x = ((\text{linepath } a \ b)) \ (2*x)$
 $x \in \{1/2..1\} \implies p \ x = ((\text{linepath } b \ c)) \ (2*x - 1)$
proof –
fix $x :: \text{real}$
assume $x \in \{0..1/2\}$
thus $p \ x = ((\text{linepath } a \ b)) \ (2*x)$
unfolding *assms*
using *make-polygonal-path.simps(4)[of a b c Nil]* **unfolding** *joinpaths-def* **by**
presburger
next
fix $x :: \text{real}$
assume $*$: $x \in \{1/2..1\}$
{ **assume** $x > 1/2$
then have $p \ x = ((\text{linepath } b \ c)) \ (2*x - 1)$
unfolding *assms*
using *make-polygonal-path.simps(4)[of a b c Nil]* **unfolding** *joinpaths-def* **by**
force
} **moreover**
{ **assume** $x = 1/2$
then have $p \ x = b \wedge ((\text{linepath } b \ c)) \ (2*x - 1) = b$
unfolding *assms*
using *make-polygonal-path.simps(4)[of a b c Nil]* **unfolding** *joinpaths-def*
by (*simp add: linepath-def mult.commute*)
}
ultimately show $p \ x = ((\text{linepath } b \ c)) \ (2*x - 1)$ **using** $*$ **by** *fastforce*
qed

lemma *polygon-linepath-images1*:

fixes $n :: \text{nat}$
assumes $n \geq 3$
assumes $\text{length } ell = n$

```

assumes  $x \in \{0..1/2\}$ 
shows  $\text{make-polygonal-path } ell \ x = ((\text{linepath } (ell \ 0) \ (ell \ 1))) \ (2*x)$ 
proof –
  have  $\text{make-polygonal-path } ell = \text{linepath } (ell \ 0) \ (ell \ 1) \ +++ \ \text{make-polygonal-path}$ 
   $(\text{drop } 1 \ ell)$ 
    using  $\text{make-polygonal-path.simps}$ 
    by  $(\text{smt } (\text{verit}, \ \text{del-insts}) \ \text{numeral-3-eq-3} \ \text{Cons-nth-drop-Suc} \ \text{One-nat-def} \ \text{Suc-1}$ 
   $\text{Suc-eq-plus1} \ \text{add-Suc-right} \ \text{assms}(1) \ \text{assms}(2) \ \text{drop0} \ \text{length-greater-0-conv} \ \text{less-add-Suc2}$ 
   $\text{list.size}(3) \ \text{not-numeral-le-zero} \ \text{nth-Cons-0} \ \text{numeral-Bit0} \ \text{order-less-le-trans} \ \text{plus-1-eq-Suc})$ 
    then show  $?thesis$ 
    using  $\text{assms } \text{make-polygonal-path.simps}$ 
    by  $(\text{simp } \text{add: } \text{joinpaths-def})$ 
qed

```

```

lemma  $\text{sum-insert} \ [\text{simp}]$ :
  assumes  $x \notin F$  and  $\text{finite } F$ 
  shows  $(\sum y \in \text{insert } x \ F. \ P \ y) = (\sum y \in F. \ P \ y) + P \ x$ 
  using  $\text{assms } \text{insert-def}$  by  $(\text{simp } \text{add: } \text{add.commute})$ 

```

```

lemma  $\text{sum-of-index-diff} \ [\text{simp}]$ :
  fixes  $f :: \text{nat} \Rightarrow 'a :: \text{comm-monoid-add}$ 
  shows  $(\sum i \in \{a..<a+b\}. \ f(i-a)) = (\sum i \in \{..<b\}. \ f(i))$ 
proof  $(\text{induction } b)$ 
  case  $0$ 
    then show  $?case$  by  $\text{simp}$ 
next
  case  $(\text{Suc } b)$ 
    then show  $?case$  by  $\text{simp}$ 
qed

```

```

lemma  $\text{sum-of-index-diff2} \ [\text{simp}]$ :
  fixes  $f :: \text{nat} \Rightarrow 'a :: \text{comm-monoid-add}$ 
  shows  $(\sum i \in \{a+c..b+c\}. \ f(i)) = (\sum i \in \{a..b\}. \ f(i+c))$ 
  using  $\text{Set-Interval.comm-monoid-add-class.sum.shift-bounds-cl-nat-ivl}$  by  $\text{blast}$ 

```

```

lemma  $\text{sum-split} \ [\text{simp}]$ :
  fixes  $f :: \text{nat} \Rightarrow 'a :: \text{comm-monoid-add}$ 
  assumes  $c \in \{a..b\}$ 
  shows  $(\sum i \in \{a..b\}. \ f \ i) = (\sum i \in \{a..c\}. \ f \ i) + (\sum i \in \{c+1..b\}. \ f \ i)$ 
  by  $(\text{metis } \text{Suc-eq-plus1} \ \text{Suc-le-mono} \ \text{assms } \text{atLeastAtMost-iff} \ \text{atLeastLessThanSuc-atLeastAtMost}$ 
   $\text{le-SucI} \ \text{sum.atLeastLessThan-concat})$ 

```

```

lemma  $\text{summation-helper}$ :
  fixes  $x :: \text{real}$ 
  fixes  $k :: \text{nat}$ 
  assumes  $1 \leq k$ 
  shows  $(2::\text{real}) * (\sum i = 1..k. \ 1 / 2^i) - 1 = (\sum i = 1..(k-1). \ (1 / (2^i)))$ 

```

proof–

have *frac-cancel*: $\forall i::\text{nat} \geq 1. 2 / (2^{\wedge}i) = 2 / (2 * (2::\text{real})^{\wedge}(i-1))$
using *power.simps(2)*[of $2::\text{real}$] **by** (*metis Suc-diff-le diff-Suc-1*)
have $(2::\text{real}) * (\sum i = 1..k. 1 / 2^{\wedge}i) = (\sum i = 1..k. (2 / 2^{\wedge}i))$
by (*simp add: sum-distrib-left*)
also have $\dots = (\sum i = 1..k. (2 / (2 * 2^{\wedge}(i-1))))$ **using** *frac-cancel* **by** *simp*
also have $\dots = (\sum i = 1..k. (1 / (2^{\wedge}(i-1))))$ **by** *force*
also have $\dots = (\sum i = 1..<(k+1). (1 / (2^{\wedge}(i-1))))$
using *Suc-eq-plus1 atLeastLessThanSuc-atLeastAtMost* **by** *presburger*
also have $\dots = (\sum i \in \{..<k\}. (1 / (2^{\wedge}i)))$
using *sum-of-index-diff*[of $\lambda i. (1 / 2^{\wedge}i) 1 k$] **by** *simp*
finally have $(2::\text{real}) * (\sum i = 1..k. 1 / 2^{\wedge}i) = (\sum i = 0..(k-1). (1 / (2^{\wedge}i)))$
by (*metis assms atLeast0AtMost diff-Suc-1 lessThan-Suc-atMost nat-le-iff-add plus-1-eq-Suc*)
then have $(2::\text{real}) * (\sum i = 1..k. 1 / 2^{\wedge}i) - 1 = (\sum i = 0..(k-1). (1 / (2^{\wedge}i))) - 1$
by *auto*
also have $\dots = (\sum i = 1..(k-1). (1 / (2^{\wedge}i))) + (1/2^{\wedge}0) - 1$
using *sum-insert*[of $0 \{1..k-1\}$ *power (1/2)*]
by (*simp add: Icc-eq-insert-lb-nat add commute*)
also have $\dots = (\sum i = 1..(k-1). (1 / (2^{\wedge}i)))$ **by** *force*
finally show $(2::\text{real}) * (\sum i = 1..k. 1 / 2^{\wedge}i) - 1 = (\sum i = 1..(k-1). (1 / (2^{\wedge}i)))$.
qed

lemma *polygon-linepath-images2*:

fixes $n k::\text{nat}$
fixes $ell::(\text{real}^{\wedge}2)$ *list*
fixes $f::\text{nat} \Rightarrow \text{real} \Rightarrow \text{real}$
assumes $n \geq 3$
assumes $0 \leq k \wedge k \leq n - 3$
assumes $\text{length } ell = n$
assumes $p: p = \text{make-polygonal-path } ell$
assumes $f = (\lambda k x. (x - (\sum i \in \{1..k\}. 1/(2^{\wedge}i))) * (2^{\wedge}(k+1)))$
assumes $x \in \{(\sum i \in \{1..k\}. 1/(2^{\wedge}i))..(\sum i \in \{1..(k+1)\}. 1/(2^{\wedge}i))\}$
shows $p x = ((\text{linepath } (ell ! k) (ell ! (k+1)) (f k x)))$
using *assms*
proof (*induct n arbitrary: ell k x p*)
case 0
then show *?case* **by** *auto*
next
case (*Suc n*)
{ **assume** $*$: $k = 0$
have $x: x \in \{0..1/2\}$ **using** $*$ *Suc.prem(6)* **by** *simp*
moreover have $f k x = 2*x$ **using** $*$ *Suc.prem(5)* **by** *simp*
ultimately have *?case*
using *polygon-linepath-images1*[of *Suc n ell x, OF Suc.prem(1) Suc.prem(3)*
 x] $*$
by (*simp add: Suc.prem(4)*)

```

} moreover
{ assume *: k ≥ 1
  then have suc-n: Suc n > 3 using Suc.prem(2) by linarith
  then have ell-is: ell = (ell!0) # (ell!1) # (ell!2) # (drop 3 ell)
    using Suc.prem(3)
    by (metis Cons-nth-drop-Suc One-nat-def Suc-1 Suc-le-lessD drop0 nat-less-le
numeral-3-eq-3)
  then have ell'-is: drop 1 ell = (ell!1) # (ell!2) # (drop 3 ell)
    by (metis One-nat-def diff-Suc-1 drop0 drop-Cons-numeral numerals(1))
  let ?ell' = drop 1 ell
  have len-ell': length ?ell' > 2 using suc-n Suc.prem(3) by simp
  let ?p' = make-polygonal-path ?ell'
  have p-tl: p = (linepath (ell ! 0) (ell ! 1)) +++ make-polygonal-path (drop 1
ell)
    using Suc.prem(4) Suc.prem(3) * make-polygonal-path.simps ell-is ell'-is
    by metis

have (∑ i = 1..k. 1 / (2 ^ i::real)) ≥ (∑ i = 1..1. 1 / (2 ^ i::real))
  using Suc.prem(2) *
proof (induct k)
  case 0
  then show ?case by auto
next
  case (Suc k)
  { assume *: 1 = Suc k
    then have ?case by auto
  } moreover { assume *: 1 < Suc k
    then have 1 ≤ k ∧ k ≤ Suc n - 3
      using Suc.prem by auto
    then have ind-h: (∑ i = 1..1. 1 / (2 ^ i::real)) ≤ (∑ i = 1..k. 1 / 2 ^ i)
      using Suc.hyps Suc.prem(2) by blast
    have (∑ i = 1..Suc k. 1 / (2 ^ i::real)) = 1/(2^(Suc k)) + (∑ i = 1..k. 1
/ (2 ^ i::real))
      using * by simp
    then have (∑ i = 1..Suc k. 1 / (2 ^ i::real)) > (∑ i = 1..k. 1 / (2 ^
i::real))
      by simp
    then have ?case using ind-h by linarith
  }
ultimately show ?case by linarith
qed
then have (∑ i = 1..k. 1 / (2 ^ i::real)) ≥ 1/2
  by auto
then have x-gteq: x ≥ 1/2 using Suc.prem(2,6)
  by (meson atLeastAtMost-iff order-trans)
have xonehalf: p x = ?p' (2*x - 1) if x-is: x = 1/2 using p-tl joinpaths-def
proof -
  have p x = (linepath (ell ! 0) (ell ! 1)) 1
    using p-tl joinpaths-def x-is

```

```

    by (metis mult.commute nle-le nonzero-divide-eq-eq zero-neq-numeral)
  then have p x = ell ! 1
    using polygon-pathfinish[of [(ell ! 0), (ell ! 1)]] unfolding pathfinish-def
    using make-polygonal-path.simps by simp
  then have p x = make-polygonal-path (drop 1 ell) 0
    using polygon-pathstart[of drop 1 ell] * len-ell' unfolding pathstart-def
    by simp
  then show ?thesis using x-is by force
qed
have x-gtonehalf: x > 1/2  $\implies$  p x = ?p' (2*x - 1) using p-tl joinpaths-def
by (smt (verit, ccfv-threshold))
then have px: p x = ?p' (2*x - 1) using xonehalf x-gtonehalf x-gteq
by linarith
{ assume k-eq: k = 1
  then have f k x = (x - ( $\sum i = 1..1. 1 / 2^i$ )) * 2^2
    using Suc.prem5 by auto
  then have f k x = 4*x - 2
    by auto
  have x  $\in$  {1/2..3/4}
    using k-eq Suc.prem6 by auto
  then have 2*x - 1  $\in$  {0..1/2} by simp
  then have ?p' (2*x - 1) = (linepath (?ell!0) (?ell!1)) (4*x - 2)
    using Suc.hyps[of k ?ell' ?p' 2*x - 1] Suc.prem5
    by (smt (verit, ccfv-SIG) suc-n diff-Suc-1 leD le-Suc-eq length-drop polygon-linepath-images1)
  also have ... = (linepath (ell!1) (ell!2)) (4*x - 2)
    using * Suc.prem3
    using ell'-is by fastforce
  also have ... = ((linepath (ell ! k) (ell ! (k+1)) (f k x))) using k-eq
    Suc.prem5 f k x
    by (smt (verit, del-insts) nat-1-add-1)
  finally have ?case using px by simp
} moreover
{ assume k-gt: k > 1
  then have f k minus: f (k-1) (2 * x - 1) = ((2 * x - 1) - ( $\sum i = 1..(k-1). 1 / 2^i$ )) * 2^k
    using Suc.prem5 by force
  have f k: f k x = (x - ( $\sum i = 1..k. 1 / 2^i$ )) * 2^(k+1)
    using Suc.prem5 by blast
  have f-is: f (k-1) (2 * x - 1) = f k x
  proof-
    have i:  $\forall i::nat \in \{2..k\}. i - 2 + 2 = i$ 
      by auto
    have f (k-1) (2 * x - 1) = (2 * x - 1 - ( $\sum i = 1..k-1. 1 / 2^i$ ))
      * 2^(k-1+1)
      unfolding Suc.prem5 by auto
    also have ... = (x - 1/2 - ( $\sum i = 1..k-1. 1 / 2^i$ ) / 2) * 2^(k+1)
      using k-gt by fastforce
    also have ... = (x - 1/2 - ( $\sum i = 1..k-1. (1 / 2^i) / 2$ )) * 2^(k+1)

```

by (*simp add: sum-divide-distrib*)
 also have ... = $(x - 1/2 - (\sum i = 1..k - 1. (1 / 2)^{\wedge} i * 1/2)) * 2^{\wedge} (k + 1)$
 by (*simp add: power-divide*)
 also have ... = $(x - 1/2 - (\sum i = 1..k - 1. (1 / 2)^{\wedge} (i+1))) * 2^{\wedge} (k + 1)$ by *force*
 also have ... = $(x - 1/2 - (\sum i = 1..<1 + (k - 1). (1 / 2)^{\wedge} (i+1))) * 2^{\wedge} (k + 1)$
 using *Suc-eq-plus1-left atLeastLessThanSuc-atLeastAtMost* by *presburger*
 also have ... = $(x - 1/2 - (\sum i = 1..<1 + (k - 1). (1 / 2)^{\wedge} (i - 1 + 2))) * 2^{\wedge} (k + 1)$
 by *auto*
 also have ... = $(x - 1/2 - (\sum i \in \{..<k - 1\}. ((1 / 2)^{\wedge} (i+2)))) * 2^{\wedge} (k + 1)$
 using *sum-of-index-diff[of ($\lambda x. (1/2)^{\wedge} (x+2)$) 1 k-1]* by *metis*
 also have ... = $(x - 1/2 - (\sum i \in \{2..<k - 1 + 2\}. ((1 / 2)^{\wedge} (i - 2 + 2)))) * 2^{\wedge} (k + 1)$
 using *sum-of-index-diff[of ($\lambda x. (1/2)^{\wedge} (x+2)$) 2 k-1]* by (*smt (verit) add.commute*)
 also have ... = $(x - 1/2 - (\sum i \in \{2..k\}. ((1 / 2)^{\wedge} (i - 2 + 2)))) * 2^{\wedge} (k + 1)$
 using *k-gt atLeastLessThanSuc-atLeastAtMost* by *force*
 also have ... = $(x - 1/2 - (\sum i \in \{2..k\}. ((1 / 2)^{\wedge} (i)))) * 2^{\wedge} (k + 1)$
 using *i* by *force*
 also have ... = $(x - (1/2 + (\sum i \in \{2..k\}. ((1 / 2)^{\wedge} (i)))) * 2^{\wedge} (k + 1)$
 by *argo*
 also have ... = $(x - (\sum i = 1..k. (1 / 2)^{\wedge} (i))) * 2^{\wedge} (k + 1)$
 using *sum-insert[of 1 {2..k} $\lambda x. (1/2)^{\wedge} x$]*
 by (*smt (verit, ccfv-SIG) Suc-1 Suc-n-not-le-n atLeastAtMost-iff atLeast-AtMost-insertL finite-atLeastAtMost k-gt less-imp-le-nat power-one-right*)
 also have ... = $(x - (\sum i = 1..k. 1 / (2^{\wedge} i))) * 2^{\wedge} (k + 1)$ by (*meson power-one-over*)
 also have ... = *f k x* using *fk* by *argo*
 finally show *?thesis* .
 qed

have *ih1*: $3 \leq n$ using *suc-n* by *force*
 have *ih2*: $0 \leq k - 1 \wedge k - 1 \leq n - 3$ using *k-gt Suc.prem(2) Suc.prem(3)*
 by *auto*
 have *ih3*: *length ?ell' = n* using *Suc.prem(3)* by *auto*
 have *ih4*: *?p' = make-polygonal-path ?ell'* by *blast*

have $2*x - 1 \geq (\sum i \in \{1..k-1\}. 1/(2^{\wedge} i))$
 proof-
 have $(2::real) * (\sum i = 1..k. 1 / 2^{\wedge} i) - 1 = (\sum i = 1..(k-1). (1 / (2^{\wedge} i)))$
 using *summation-helper k-gt* by *auto*
 moreover have $x \geq (\sum i = 1..k. 1 / 2^{\wedge} i)$ using *Suc.prem(6)* by *presburger*

```

    ultimately show  $2*x - 1 \geq (\sum i \in \{1..k-1\}. 1/(2^i))$  by linarith
  qed
  moreover have  $2*x - 1 \leq (\sum i \in \{1..k\}. 1/(2^i))$ 
  proof-
    have  $(2::real) * (\sum i \in \{1..(k+1)\}. 1/(2^i)) - 1 = (\sum i \in \{1..k\}. 1/(2^i))$ 
    using summation-helper[of k + 1] k-gt by auto
    moreover have  $x \leq (\sum i \in \{1..(k+1)\}. 1/(2^i))$  using Suc.premis(6)
  by presburger
    ultimately show ?thesis by linarith
  qed
    ultimately have  $2*x - 1 \in \{(\sum i \in \{1..k-1\}. 1/(2^i))..(\sum i \in \{1..k\}. 1/(2^i))\}$  by presburger
    then have ih5:  $2*x - 1 \in \{(\sum i \in \{1..k-1\}. 1/(2^i))..(\sum i \in \{1..k-1+1\}. 1/(2^i))\}$ 
    using k-gt by auto

    have  $p = \text{make-polygonal-path } (ell!0 \# ell!1 \# ell!2 \# (\text{drop } 3 \text{ ell}))$ 
    using ell-is Suc.premis(4) by argo
    then have  $p = (\text{linepath } (ell!0) (ell!1)) \text{+++ make-polygonal-path } (ell!1 \# ell!2 \# (\text{drop } 3 \text{ ell}))$ 
    using make-polygonal-path.simps by auto
    then have  $p \ x = ?p' (2*x - 1)$  unfolding joinpaths-def using x-gteq px by fastforce
    also have  $\dots = (\text{linepath } (?ell!(k-1)) (?ell!k)) (f (k-1) (2*x - 1))$ 
    using Suc.hyps[OF ih1 ih2 ih3 ih4 Suc.premis(5), of 2*x - 1, OF ih5] using k-gt by auto
    also have  $\dots = (\text{linepath } (ell!k) (ell!(k+1))) (f (k-1) (2*x - 1))$ 
    using Suc.premis(2) Suc.premis(3)
    by (smt (verit, del-insts) add-implies-diff ell'-is ell-is k-gt nth-Cons-pos order-le-less-trans trans-less-add1 zero-less-one-class.zero-le-one)
    also have  $\dots = (\text{linepath } (ell!k) (ell!(k+1))) (f k x)$ 
    using f-is by auto
    finally have ?case .
  }
  ultimately have ?case using Suc.premis(2) * by linarith
}
ultimately show ?case
using Suc.premis by linarith
qed

```

lemma *polygon-linepath-images3*:

```

  fixes  $n k:: nat$ 
  fixes  $ell:: (real^2) \text{ list}$ 
  assumes  $n \geq 3$ 
  assumes  $\text{length } ell = n$ 
  assumes  $p = \text{make-polygonal-path } ell$ 
  assumes  $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$ 
  assumes  $f = (\lambda x. (x - (\sum i \in \{1..n-2\}. 1/(2^i))) * (2^{(n-2)}))$ 

```

```

shows  $p x = (\text{linepath } (\text{ell} ! (n-2)) (\text{ell} ! (n-1))) (f x)$ 
using assms
proof (induct n arbitrary: ell k x p f)
  case 0
  then show ?case by auto
next
case (Suc n)
  { assume *:  $\text{Suc } n = 3$ 
    then have ell-is:  $\text{ell} = [\text{ell} ! 0, \text{ell} ! 1, \text{ell} ! 2]$ 
      using Suc.prem(2)
      by (metis Cons-nth-drop-Suc One-nat-def Suc-1 cancel-comm-monoid-add-class.diff-cancel
drop0 length-0-conv length-drop lessI less-add-Suc2 numeral-3-eq-3 plus-1-eq-Suc
zero-less-Suc)
      have  $(\sum i = 1..(\text{Suc } n)-2. 1 / ((2 \wedge i)::\text{real})) = (\sum i \in \{1\}. 1 / ((2 \wedge i)::\text{real}))$ 
        by (simp add: *)
      then have eq1:  $(\sum i = 1..(\text{Suc } n)-2. 1 / ((2 \wedge i)::\text{real})) = 1/2$ 
        by auto
      then have f-is:  $f = (\lambda x. (x - (1/2)) * 2)$  using * Suc.prem(5) by auto
      have  $x \in \{(1/2)::\text{real}..1\}$  using eq1 Suc.prem(4) by metis
      moreover then have  $p x = \text{linepath } (\text{ell} ! 1) (\text{ell} ! 2) (2 * x - 1)$ 
        using triangle-linepath-images(2) using ell-is Suc.prem(3) by blast
      moreover have  $f x = 2*x - 1$  using f-is by simp
      ultimately have  $p x = (\text{linepath } (\text{ell} ! ((\text{Suc } n)-2)) (\text{ell} ! ((\text{Suc } n)-1))) (f x)$ 
        using * Suc.prem ell-is
        by (metis One-nat-def Suc-1 diff-Suc-1 diff-Suc-Suc numeral-3-eq-3)
    } moreover
  { assume *:  $\text{Suc } n > 3$ 
    let ?ell' = drop 1 ell
    let ?p' = make-polygonal-path ?ell'
    let ?x' =  $2*x - 1$ 
    let ?f' =  $(\lambda x. (x - (\sum i \in \{1..n-2\}. 1/(2 \wedge i))) * (2 \wedge (n-2)))$ 
    have ell-is:  $\text{ell} = \text{ell} ! 0 \# \text{ell} ! 1 \# \text{ell} ! 2 \# (\text{drop } 3 \text{ ell})$ 
      by (metis * Cons-nth-drop-Suc One-nat-def Suc.prem(2) Suc-1 drop0 le-Suc-eq
linorder-not-less numeral-3-eq-3 zero-less-Suc)
      then have p-tl:  $p = (\text{linepath } (\text{ell} ! 0) (\text{ell} ! 1)) \text{ +++ } \text{make-polygonal-path}$ 
        (drop 1 ell)
        using make-polygonal-path.simps(4)[of ell!0 ell!1 ell!2 drop 3 ell]
        by (metis One-nat-def Suc.prem(3) drop-0 drop-Suc-Cons)
      have sum-split:  $(\sum i = 1..\text{Suc } n - 2. 1 / (2 \wedge i::\text{real})) = 1/(2 \wedge 1::\text{real}) + (\sum i$ 
        =  $2..\text{Suc } n - 2. 1 / (2 \wedge i::\text{real}))$ 
        using *
      by (metis Suc-1 Suc-eq-plus1 Suc-lessD add-le-imp-le-diff diff-Suc-Suc eval-nat-numeral(3)
less-Suc-eq-le sum.atLeast-Suc-atMost)
      let ?k =  $\text{Suc } n$ 
      have helper-arith:  $\bigwedge i. i > 0 \implies 1 / (2 \wedge i::\text{real}) > 0$  by simp
      have  $k \geq 2 \implies (\sum i = 2..k. 1 / (2 \wedge i::\text{real})) > 0$  for k
      proof (induct k)
        case 0
        then show ?case by auto

```

```

next
  case (Suc k)
  {assume *: Suc k = 2
   then have ( $\sum i = 2..Suc\ k. 1 / (2 \wedge i::real)$ ) = ( $\sum i = 2..2. 1 / (2 \wedge i::real)$ )
   by presburger
   then have ?case
   using helper-arith
   by (simp add: *)
  } moreover {assume *: Suc k > 2
  then have ind-h:  $0 < (\sum i = 2..k. 1 / (2 \wedge i::real))$ 
  using Suc.hyps less-Suc-eq-le by blast
  have ( $\sum i = 2..Suc\ k. 1 / (2 \wedge i::real)$ ) = ( $\sum i = 2..k. 1 / (2 \wedge i::real)$ )
+  $1 / (2 \wedge (Suc\ k)::real)$ 
  using Suc.prem1 add.commute by auto
  then have ?case using ind-h helper-arith
  by (smt (verit) divide-less-0-1-iff zero-le-power)
  }
ultimately show ?case
using Suc.prem1 by linarith
qed
then have ( $\sum i = 2..Suc\ n - 2. 1 / (2 \wedge i::real)$ ) > 0
using * by auto
then have ( $\sum i = 1..Suc\ n - 2. 1 / (2 \wedge i::real)$ ) > 1/2
using sum-split by auto
then have  $x > 1/2$  using Suc.prem1(4)
by (smt (verit, del-insts) atLeastAtMost-iff linorder-not-le order-le-less-trans)
then have  $p'x'-eq-px: ?p' ?x' = p\ x$  unfolding joinpaths-def by (simp add:
joinpaths-def p-tl)

have 1:  $n \geq 3$  using * by auto
have 2:  $length\ ?ell' = n$  using Suc.prem1(2) by simp
have 3:  $?p' = make\_polygonal\_path\ ?ell'$  by auto
have  $x \leq 1$  using Suc.prem1(4) by auto
then have  $x'-lteq: 2*x - 1 \leq 1$  by auto
have  $x \geq (\sum i = 1..Suc\ n - 2. 1 / 2 \wedge i)$ 
using Suc.prem1(4) by auto
then have  $x'-gteq: ?x' \geq (\sum i = 1..n - 2. 1 / 2 \wedge i)$ 
using summation-helper[of Suc n - 2] *
by (smt (verit) Suc.prem1(1) Suc-1 Suc-diff-le Suc-leD Suc-le-mono diff-Suc-1
diff-Suc-eq-diff-pred eval-nat-numeral(3))
have 4:  $?x' \in \{(\sum i = 1..n - 2. 1 / 2 \wedge i)..1\}$  using Suc.prem1(4)
using summation-helper[of Suc n - 2] *  $x'-lteq\ x'-gteq\ atLeastAtMost-iff$  by
blast
have 5:  $?f' = (\lambda x. (x - (\sum i = 1..n - 2. 1 / 2 \wedge i)) * 2 \wedge (n - 2))$  by auto
have  $f\ x = (x - (\sum i = 1..Suc\ n - 2. 1 / 2 \wedge i)) * 2 \wedge (n - 2)*2$ 
proof -
have  $(\lambda r. (r - (\sum n = 1..n - 1. 1 / 2 \wedge n)) * 2 \wedge (n - 1)) = f$ 
by (simp add: Suc.prem1(5))

```

```

then have  $2^{n-1} * (x - (\sum_{n=1..n-1} 1 / 2^n)) = f x$ 
using Groups.mult-ac(2) by blast
then have  $(x - (\sum_{n=1..n-1} 1 / 2^n)) * (2^{n-1} * 2) = f x$ 
by (metis (no-types) Groups.mult-ac(2) Suc.premis(2) diff-Suc-1 diff-Suc-Suc
ell-is length-Cons power.simps(2))
then show ?thesis
by (metis (no-types) Groups.mult-ac(1) Suc-1 diff-Suc-Suc)
qed
then have fx-is:  $f x = (2*x - 2*(\sum_{i=1..Suc\ n-2} 1 / 2^i)) * 2^{n-2}$ 
by argo
have sum-is:  $1 + (\sum_{i=1..n-2} 1 / (2^{i::real})) = 2*(\sum_{i=1..Suc\ n-2} 1 / (2^{i::real}))$ 
proof -
have sum-ish1:  $(\sum_{i=1..Suc\ n-2} 1 / (2^{i::real})) = 1/2 + (\sum_{i=2..Suc\ n-2} 1 / (2^{i::real}))$ 
by (metis power-one-right sum-split)
have  $n \geq 2 \implies 2*(\sum_{i=2..n-1} 1 / (2^{i::real})) = (\sum_{i=1..n-2} 1 / (2^{i::real}))$ 
proof (induct n)
case 0
then show ?case by auto
next
case (Suc n)
{assume *: Suc n = 2
then have ?case by auto
} moreover {assume *: Suc n > 2
then have ind-h:  $2 * (\sum_{i=2..n-1} 1 / (2^{i::real})) = (\sum_{i=1..n-2} 1 / (2^{i::real}))$ 
using Suc by fastforce
have mult:  $2*1/(2^{Suc\ n-1}::real) = 1/(2^{n-1}::real)$ 
using *
by (smt (z3) One-nat-def add-diff-inverse-nat bot-nat-0.not-eq-extremum
diff-Suc-1 div-by-1 le-zero-eq less-Suc-eq-le mult commute nonzero-mult-div-cancel-left
nonzero-mult-divide-mult-cancel-left plus-1-eq-Suc power-Suc zero-less-numeral)
have sum-prop:  $\bigwedge a::nat. \bigwedge f::nat \Rightarrow real. (\sum_{i=1..a} (f\ i)) + (f\ (a+1)) = (\sum_{i=1..a+1} (f\ i))$ 
by auto
have  $n - 2 + 1 = n - 1$ 
using * by auto
then have sum-same:  $(\sum_{i=1..n-2} 1 / (2^{i::real})) + 1 / 2^{n-1} = (\sum_{i=1..n-1} 1 / (2^{i::real}))$ 
using * sum-prop[of  $\lambda i. 1 / (2^{i::real})$  n-2] by metis
have  $2 * (\sum_{i=2..Suc\ n-1} 1 / (2^{i::real})) = 2 * ((\sum_{i=2..n-1} 1 / (2^{i::real})) + 1/(2^{Suc\ n-1}::real))$ 
using *
by (smt (z3) add-2-eq-Suc add-diff-inverse-nat diff-Suc-1 distrib-left-numeral
ind-h not-less-eq sum.cl-ivl-Suc)
then have  $2 * (\sum_{i=2..Suc\ n-1} 1 / (2^{i::real})) = (\sum_{i=1..n-1} 1 / (2^{i::real}))$ 

```

```

2. 1 / (2 ^ i::real)) + 2*1/(2^(Suc n - 1)::real)
  using ind-h by argo
  then have 2 * (∑ i = 2..Suc n - 1. 1 / (2 ^ i::real)) = (∑ i = 1..n -
2. 1 / (2 ^ i::real)) + 1/(2^(n - 1)::real)
  using * mult by auto
  then have ?case using sum-same by auto
}
ultimately show ?case by fastforce
qed
then have sum-ish2:2*(∑ i = 2..Suc n - 2. 1 / (2 ^ i::real)) = (∑ i =
1..n - 2. 1 / (2 ^ i::real))
  using * by auto
  show ?thesis using sum-ish1 sum-ish2 by simp
qed
have ?p' ?x' = (linepath (?ell' ! (n-2)) (?ell' ! (n-1))) (?f' ?x')
  using Suc.hyps[OF 1 2 3 4 5] by blast
moreover have ?f' ?x' = f x
  using Suc.prem5 fx-is sum-is
  by (smt (verit, best))
moreover have ?ell' ! (n-2) = ell ! ((Suc n)-2)
  by (metis Nat.diff-add-assoc One-nat-def Suc.prem1 Suc.prem2 Suc-1
add-diff-cancel-left le-add1 nth-drop numeral-3-eq-3 plus-1-eq-Suc)
moreover have ?ell' ! (n-1) = ell ! ((Suc n)-1)
  using Suc.prem1 Suc.prem2 by auto
ultimately have ?case using p'x'-eq-px by presburger
}
ultimately show ?case using Suc.prem1 by linarith
qed

```

8 A Triangle is a Polygon

lemma *not-collinear-linepaths-intersect-helper*:

```

assumes not-collinear: ¬collinear {a,b,c}
assumes 0 ≤ k1
assumes k1 ≤ 1
assumes 0 ≤ k2
assumes k2 ≤ 1
assumes eo: k2 = 0 ⇒ k1 ≠ 1
shows ¬((linepath a b) k1 = (linepath b c) k2)

```

proof –

```

have a-neq-b: a ≠ b
  using not-collinear
  by auto
then have nonz-1: a - b ≠ 0
  by auto
have b-neq-c: b ≠ c
  using not-collinear
  by auto
then have nonz-2: b - c ≠ 0

```

```

    by auto
  have  $\neg$  collinear {a-b, 0, c-b}
    using not-collinear
    by (metis NO-MATCH-def collinear-3 insert-commute)
  then have notcollinear:  $\neg$  collinear {0, a-b, c-b}
    by (simp add: insert-commute)
  have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies (a - k1 *_R a) + k1 *_R b = (b - k2 *_R b) + k2 *_R c$ 
    by (metis add-diff-cancel scaleR-collapse)
  then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies (1 - k1) *_R a + k1 *_R b - b = -k2 *_R b + k2 *_R c$ 
    by (metis (no-types, lifting) add-diff-cancel-left scaleR-collapse scaleR-minus-left uminus-add-conv-diff)
  then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies (1 - k1) *_R a + k1 *_R b - b = k2 *_R (c-b)$ 
    by (simp add: scaleR-right-diff-distrib)
  then have rewrite:  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies (1-k1)*_R(a - b) = k2 *_R (c-b)$ 
    by (metis add-diff-cancel-right scaleR-collapse scaleR-right-diff-distrib)
  {assume *:  $k2 \neq 0$ 
    then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies c - b = ((1-k1)/k2)*_R(a - b)$ 
      using rewrite assms(2-3)
      by (smt (verit, ccfv-SIG) vector-fraction-eq-iff)
    then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies$  collinear {0, a-b, c-b}
      using collinear-lemma[of a -b c-b] by auto
    then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies$  False
      using notcollinear by auto
  } moreover {assume *:  $k2 = 0$ 
    then have  $k1 \neq 1$ 
      using assms by auto
    then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies a - b = (k2/(1-k1)) *_R (c-b)$ 
      using rewrite
      by (smt (verit, ccfv-SIG) vector-fraction-eq-iff)
    then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies$  collinear {0, a-b, c-b}
      using collinear-lemma[of c-b a-b]
      by (simp add: insert-commute)
    then have  $(1 - k1) *_R a + k1 *_R b = (1 - k2) *_R b + k2 *_R c \implies$  False
      using notcollinear by auto
  }
  ultimately show ?thesis
    unfolding linpath-def
    by blast
qed

```

```

lemma not-collinear-linepaths-intersect-helper-2:
  assumes not-collinear:  $\neg \text{collinear } \{a,b,c\}$ 
  assumes  $0 \leq k1$ 
  assumes  $k1 \leq 1$ 
  assumes  $0 \leq k2$ 
  assumes  $k2 \leq 1$ 
  assumes eo:  $k1 = 0 \implies k2 \neq 1$ 
  shows  $\neg ((\text{linepath } a \ b) \ k1 = (\text{linepath } c \ a) \ k2)$ 
  using not-collinear-linepaths-intersect-helper[of c a b k2 k1] assms
  by (simp add: insert-commute)

lemma not-collinear-loopfree-path:  $\bigwedge a \ b \ c::\text{real}^2. \neg \text{collinear } \{a,b,c\} \implies \text{loop-free}$ 
  ( $(\text{linepath } a \ b) \ +++ \ (\text{linepath } b \ c)$ )
proof -
  fix  $a \ b \ c::\text{real}^2$ 
  assume not-collinear:  $\neg \text{collinear } \{a,b,c\}$ 
  then have a-neq-b:  $a \neq b$ 
    by auto
  have b-neq-c:  $b \neq c$ 
    using not-collinear
    by auto
  have  $\bigwedge x \ y::\text{real}. (\text{linepath } a \ b \ +++ \ \text{linepath } b \ c) \ x = (\text{linepath } a \ b \ +++ \ \text{linepath}$ 
   $b \ c) \ y \implies$ 
     $x < y \implies$ 
     $x = 0 \longrightarrow y \neq 1 \implies 0 \leq x \implies x \leq 1 \implies 0 \leq y \implies y \leq 1 \implies \text{False}$ 
proof -
  fix  $x \ y::\text{real}$ 
  assume same-eval:  $(\text{linepath } a \ b \ +++ \ \text{linepath } b \ c) \ x = (\text{linepath } a \ b \ +++$ 
   $\text{linepath } b \ c) \ y$ 
  assume x-neq-y:  $x < y$ 
  assume x-zero-imp:  $x = 0 \longrightarrow y \neq 1$ 
  assume x-gt:  $0 \leq x$ 
  assume x-lt:  $x \leq 1$ 
  assume y-gt:  $0 \leq y$ 
  assume y-lt:  $y \leq 1$ 
  {assume *:  $x \leq 1/2 \wedge y \leq 1/2$ 
  then have  $(1 - 2 * x) *_{\mathbb{R}} a + (2 * x) *_{\mathbb{R}} b = (1 - 2 * y) *_{\mathbb{R}} a + (2 * y)$ 
 $*_{\mathbb{R}} b \implies \text{False}$ 
  using x-gt y-gt x-neq-y a-neq-b linepath-loop-free[of a b]
  by (smt (z3) add-diff-cancel-left add-diff-cancel-right' add-diff-eq scaleR-cancel-left
  scaleR-left-diff-distrib)
  then have False
  using * same-eval unfolding joinpaths-def linepath-def
  by auto
} moreover {assume *:  $x > 1/2 \wedge y > 1/2$ 
have False
  using x-lt y-lt x-neq-y b-neq-c linepath-loop-free[of b c]
  using * same-eval unfolding joinpaths-def linepath-def
by (smt (z3) add-diff-cancel-left add-diff-cancel-right' add-diff-eq scaleR-cancel-left

```

```

scaleR-collapse scaleR-left-diff-distrib)
} moreover {assume *:  $x \leq 1/2 \wedge y > 1/2$ 

  then have lp-eq: (linepath a b) (2 * x) = (linepath b c) (2 * y - 1)
    using * same-eval unfolding joinpaths-def
    by auto
  have (2 * y - 1) = 0  $\longrightarrow$  (2*x)  $\neq$  1  $\wedge$  0  $\leq$  (2*x)  $\wedge$  (2*x)  $\leq$  1  $\wedge$  0  $\leq$  (2
* y - 1)  $\wedge$  (2 * y - 1)  $\leq$  1
    using x-lt x-gt x-neq-y * by auto
  then have False
    using lp-eq not-collinear-linepaths-intersect-helper[of a b c 2*x 2 * y - 1]
    not-collinear
    using * x-gt y-lt by auto
}
ultimately show False
  using x-lt y-lt x-neq-y
  by linarith
qed
then have  $\bigwedge x y :: \text{real}. (\text{linepath } a \ b \ \text{+++} \ \text{linepath } b \ c) \ x = (\text{linepath } a \ b \ \text{+++} \ \text{linepath } b \ c) \ y \implies$ 
   $x \neq y \implies$ 
   $x = 0 \longrightarrow y \neq 1 \implies x = 1 \longrightarrow y \neq 0 \implies 0 \leq x \implies x \leq 1 \implies 0 \leq y$ 
 $\implies y \leq 1 \implies \text{False}$ 
  by (metis linorder-less-linear)
then show loop-free (linepath a b +++ linepath b c)
  unfolding loop-free-def
  by (metis atLeastAtMost-iff)
qed

lemma triangle-is-polygon:  $\bigwedge a \ b \ c. \neg \text{collinear } \{a, b, c\} \implies \text{polygon } (\text{make-triangle } a \ b \ c)$ 
proof -
  fix a b c :: real^2
  assume not-coll:  $\neg \text{collinear } \{a, b, c\}$ 
  then have a-neq-b:  $a \neq b$ 
    by auto
  have b-neq-c:  $b \neq c$ 
    using not-coll
    by auto
  have a-neq-c:  $c \neq a$ 
    using not-coll
    using collinear-3-eq-affine-dependent by blast
  let ?vts = [a, b, c, a]
  have polygonal-path: polygonal-path (make-polygonal-path [a, b, c, a])
    by (metis Collect-const UNIV-I image-eqI polygonal-path-def)
  then have path: path (make-polygonal-path [a, b, c, a])
    by auto
  then have closed-path: closed-path (make-polygonal-path [a, b, c, a])
    unfolding closed-path-def using polygon-pathstart polygon-pathfinish

```

```

    by auto
  let ?seg1 = (linepath a b) +++ (linepath b c)
  have lf1: loop-free ((linepath a b) +++ (linepath b c))
    using not-collinear-loopfree-path not-coll
  by auto
  then have  $\forall x \in \{0..1\}. \forall y \in \{0..1\}. ?seg1\ x = ?seg1\ y \longrightarrow x = y$ 
    using a-neq-c unfolding loop-free-def
    by (metis (no-types, lifting) path-defs(2) pathfinish-def pathfinish-join pathfin-
ish-linepath pathstart-join pathstart-linepath)
  let ?seg2 = (linepath b c) +++ (linepath c a)
  have lf2: loop-free ((linepath b c) +++ (linepath c a))
    using not-collinear-loopfree-path not-coll
  by (simp add: insert-commute)
  then have  $\forall x \in \{0..1\}. \forall y \in \{0..1\}. ?seg2\ x = ?seg2\ y \longrightarrow x = y$ 
    using a-neq-b unfolding loop-free-def
    by (metis (no-types, lifting) path-defs(2) pathfinish-def pathfinish-join pathfin-
ish-linepath pathstart-join pathstart-linepath)
  let ?seg3 = (linepath c a) +++ (linepath a b)
  have lf3: loop-free ((linepath c a) +++ (linepath a b))
    using not-collinear-loopfree-path not-coll
  by (simp add: insert-commute)
  then have  $\forall x \in \{0..1\}. \forall y \in \{0..1\}. ?seg3\ x = ?seg3\ y \longrightarrow x = y$ 
    using b-neq-c unfolding loop-free-def
    by (metis (no-types, lifting) path-defs(2) pathfinish-def pathfinish-join pathfin-
ish-linepath pathstart-join pathstart-linepath)
  have mpp-is:  $\forall x \in \{0..1\}. \text{make-polygonal-path } [a, b, c, a]\ x = ((\text{linepath } a\ b)
+++ (\text{linepath } b\ c) +++ (\text{linepath } c\ a))\ x$ 
  by auto
  have x-in-int1:  $\forall x \in \{0..(1/2)\}. \text{make-polygonal-path } [a, b, c, a]\ x = ((\text{linepath }
a\ b))\ (2*x)$ 
  using mpp-is
  unfolding joinpaths-def by auto
  have x-in-int2:  $\forall x \in \{1/2 <..(3/4)\}. \text{make-polygonal-path } [a, b, c, a]\ x = ((\text{linepath }
b\ c))\ (2*(2*x - 1))$ 
  using mpp-is unfolding joinpaths-def
  by auto
  have x-in-int3:  $\forall x \in \{3/4 <..1\}. \text{make-polygonal-path } [a, b, c, a]\ x = ((\text{linepath }
c\ a))\ (2 * (2 * x - 1) - 1)$ 
  using mpp-is unfolding joinpaths-def
  by auto
  have  $\bigwedge x\ y. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x \neq y \wedge (x = 0 \longrightarrow y \neq 1) \wedge
(x = 1 \longrightarrow y \neq 0) \implies \text{make-polygonal-path } [a, b, c, a]\ x = \text{make-polygonal-path }
[a, b, c, a]\ y \implies \text{False}$ 
  proof -
    fix x y: real
    assume big:  $0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x \neq y \wedge (x = 0 \longrightarrow y \neq 1)
\wedge (x = 1 \longrightarrow y \neq 0)$ 
    assume false-hyp:  $\text{make-polygonal-path } [a, b, c, a]\ x = \text{make-polygonal-path } [a,
b, c, a]\ y$ 

```

```

{assume *: x ∈ {0..(1/2)}}
  then have x-eval: make-polygonal-path [a, b, c, a] x = ((linepath a b)) (2*x)
    using x-in-int1 by auto
  {assume **: y ∈ {0..(1/2)}}
    then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath a b))
(2*y)
      using x-in-int1 by auto
    then have ((linepath a b)) (2*x) = ((linepath a b)) (2*y)
      using false-hyp x-eval y-eval by auto
    then have False
      using linepath-loop-free big * **
      unfolding loop-free-def
        using a-neq-b add-diff-cancel-left add-diff-cancel-right' add-diff-eq
linepath-def scaleR-cancel-left scaleR-collapse scaleR-left-diff-distrib
      by (smt (verit))
    } moreover {assume **: y ∈ {(1/2)<..(3/4)}}
      then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath b c))
(2*(2*y - 1))
        using x-in-int2 by auto
      then have ((linepath a b)) (2*x) = ((linepath b c)) (2*(2*y - 1))
        using false-hyp x-eval y-eval by auto
      then have False
        using big * ** not-collinear-linepaths-intersect-helper[of a b c 2*x
(2*(2*y - 1))] not-coll
        by auto
      } moreover {assume **: y ∈ {(3/4)<..1}}
        then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath c a))
((2 * (2 * y - 1) - 1))
          using x-in-int3 by auto
        then have ((linepath a b)) (2*x) = ((linepath c a)) ((2 * (2 * y - 1)
- 1))
          using false-hyp x-eval y-eval by auto
        then have False
          using big * ** not-collinear-linepaths-intersect-helper-2[of a b c (2*x)
((2 * (2 * y - 1) - 1))] not-coll
          by auto
        }
      ultimately have False
        using big
        by (metis atLeastAtMost-iff greaterThanAtMost-iff linorder-not-le)
    } moreover {assume *: x ∈ {(1/2)<..(3/4)}}
      then have x-eval: make-polygonal-path [a, b, c, a] x = ((linepath b c))
(2*(2*x - 1))
        using x-in-int2 by auto
      {assume **: y ∈ {0..(1/2)}}
        then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath a b))
(2*y)
          using x-in-int1 by auto
        then have lp-eq: ((linepath a b)) (2*y) = ((linepath b c)) (2*(2*x - 1))

```

```

    using false-hyp x-eval y-eval by auto
    have 2 * (2 * x - 1) ≠ 0
    using * by auto
    then have False
    using lp-eq big * ** not-collinear-linepaths-intersect-helper[of a b c 2*y
(2*(2*x - 1))] not-coll
    by auto
  } moreover {assume **: y ∈ {(1/2)<..(3/4)}}
    then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath b c))
(2*(2*y - 1))
    using x-in-int2 by auto
    then have lp-eq: ((linepath b c)) (2*(2*y - 1)) = ((linepath b c))
(2*(2*x - 1))
    using false-hyp x-eval y-eval by auto
    then have False
    using linepath-loop-free[OF b-neq-c] big * **
    unfolding loop-free-def
    using add-diff-cancel-left add-diff-cancel-right' add-diff-eq linepath-def
scaleR-cancel-left scaleR-collapse scaleR-left-diff-distrib
    by (smt (verit) b-neq-c)
  } moreover {assume **: y ∈ {(3/4)<..1}}
    then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath c a))
((2 * (2 * y - 1) - 1))
    using x-in-int3 by auto
    then have lp-eq: ((linepath b c)) (2*(2*x - 1)) = ((linepath c a)) ((2
* (2 * y - 1) - 1))
    using false-hyp x-eval y-eval
    by auto
    have not-coll2: ¬ collinear {b, c, a}
    using not-coll
    by (simp add: insert-commute)
    have 2 * (2 * x - 1) ≠ 0
    using * by auto
    then have False using lp-eq
    using big * ** not-collinear-linepaths-intersect-helper[of b c a 2*(2*x
- 1) (2 * (2 * y - 1) - 1)] not-coll2
    by auto
  }
} ultimately have False
using big
by (metis atLeastAtMost-iff greaterThanAtMost-iff linorder-not-le)
} moreover {assume *: x ∈ {(3/4)<..1}}
    then have x-eval: make-polygonal-path [a, b, c, a] x = ((linepath c a)) ((2
* (2 * x - 1) - 1))
    using x-in-int3 by auto
    {assume **: y ∈ {0..(1/2)}}
    then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath a b))
(2*y)
    using x-in-int1 by auto

```

```

then have lp-eq: ((linepath c a)) ((2 * (2 * x - 1) - 1)) = ((linepath
a b)) (2*y)
  using x-eval y-eval
  using false-hyp by presburger
  have not-coll2:  $\neg$  collinear {c, a, b}
  using not-coll
  by (simp add: insert-commute)
  have ((2 * (2 * x - 1) - 1))  $\neq$  0
  using * by auto
  then have False
  using lp-eq big * ** not-coll2
  not-collinear-linepaths-intersect-helper[of c a b (2 * (2 * x - 1) - 1)
2*y]
  by auto
} moreover {assume **:  $y \in \{(1/2) <.. (3/4)\}$ 
then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath b c))
(2*(2*y - 1))
  using x-in-int2 by auto
then have lp-eq: ((linepath b c)) (2*(2*y - 1)) = ((linepath c a)) ((2
* (2 * x - 1) - 1))
  using x-eval y-eval false-hyp
  using false-hyp by presburger
  have not-coll2:  $\neg$  collinear {b, c, a}
  using not-coll
  by (simp add: insert-commute)
  have ((2 * (2 * x - 1) - 1))  $\neq$  0
  using * by auto
then have False
  using lp-eq big * ** not-coll2
  not-collinear-linepaths-intersect-helper[of b c a (2*(2*y - 1)) (2 * (2
* x - 1) - 1)]
  by auto
} moreover {assume **:  $y \in \{(3/4) <.. 1\}$ 
then have y-eval: make-polygonal-path [a, b, c, a] y = ((linepath c a))
((2 * (2 * y - 1) - 1))
  using x-in-int3 by auto
then have ((linepath c a)) ((2 * (2 * y - 1) - 1)) = ((linepath c a))
((2 * (2 * x - 1) - 1))
  using x-eval y-eval false-hyp
  using false-hyp by presburger
then have False
  using linepath-loop-free[OF a-neq-c] big * **
  unfolding loop-free-def
  using add-diff-cancel-left add-diff-cancel-right' add-diff-eq linepath-def
scaleR-cancel-left scaleR-collapse scaleR-left-diff-distrib
  by (smt (verit) a-neq-c add-diff-cancel-left')
}
ultimately have False
using big

```

```

      by (metis atLeastAtMost-iff greaterThanAtMost-iff linorder-not-le)
    }
  ultimately show False using big
    by (metis atLeastAtMost-iff greaterThanAtMost-iff linorder-not-le)
qed
then have loop-free: loop-free (make-polygonal-path [a, b, c, a])
  unfolding loop-free-def
  by (meson atLeastAtMost-iff)
show polygon (make-triangle a b c)
  unfolding make-triangle-def polygon-def simple-path-def
  using polygonal-path closed-path loop-free by auto
qed

```

```

lemma have-wraparound-vertex:
  assumes polygon p
  assumes p = make-polygonal-path vts
  shows vts = (take (length vts - 1) vts)@[vts ! 0]
proof -
  have card (set vts) ≥ 3
    using polygon-at-least-3-vertices assms by auto
  then have nonempty: vts ≠ []
    by auto
  then have vts = (take (length vts - 1) vts)@[vts ! (length vts - 1)]
    by (metis append-butlast-last-id butlast-conv-take last-conv-nth)
  then show ?thesis
    using assms(1) unfolding polygon-def closed-path-def
    using polygon-pathstart[OF nonempty assms(2)] polygon-pathfinish[OF nonempty
assms(2)]
    by presburger
qed

```

```

lemma polygon-at-least-3-vertices-wraparound:
  assumes polygon p
  assumes p = make-polygonal-path vts
  shows card (set (take (length vts - 1) vts)) ≥ 3
proof -
  let ?distinct-vts = take (length vts - 1) vts
  have card-vts: card (set vts) ≥ 3
    using polygon-at-least-3-vertices assms by auto
  then have vts-is: vts = ?distinct-vts@[vts ! 0]
    using have-wraparound-vertex assms by auto
  then have ?distinct-vts ≠ []
    using card-vts
  by (metis One-nat-def append-Nil distinct-card distinct-singleton eval-nat-numeral(3)
length-append-singleton list.size(3) not-less-eq-eq one-le-numeral)
  then have vts ! 0 ∈ set ?distinct-vts
    by (metis ‹vts = take (length vts - 1) vts @ [vts ! 0]› length-greater-0-conv)

```

```

nth-append nth-mem)
  then have card (set ?distinct-vts) = card (set vts)
    using vts-is
  by (metis Un-insert-right append.right-neutral insert-absorb list.set(2) set-append)
  then show ?thesis using card-vts by auto
qed

```

9 Polygon Vertex Rotation

definition *rotate-polygon-vertices*:: 'a list \Rightarrow nat \Rightarrow 'a list
where *rotate-polygon-vertices* ell i =
 (let ell1 = rotate i (butlast ell) in ell1 @ [ell1 ! 0])

lemma *rotate-polygon-vertices-same-set*:
assumes *polygon* (make-polygonal-path vts)
shows set (rotate-polygon-vertices vts i) = set vts

proof –

```

have card-gteq: card (set vts)  $\geq$  3
  using polygon-at-least-3-vertices assms
  by auto
then have len-gteq: length vts  $\geq$  3
  using card-length order-trans by blast
let ?ell1 = rotate i (take (length vts - 1) vts)
have inset: vts ! 0 = vts ! (length vts - 1)
  using assms polygon-pathstart polygon-pathfinish unfolding polygon-def closed-path-def
  by (metis len-gteq list.size(3) not-numeral-le-zero)
have set vts = set (take (length vts - 1) vts)  $\cup$  {vts ! (length vts - 1)}
  by (metis Cons-nth-drop-Suc One-nat-def Un-insert-right assms card.empty
diff-zero drop-rev length-greater-0-conv list.set(1) list.set(2) not-numeral-le-zero
order.refl polygon-at-least-3-vertices rev-nth set-rev sup-bot.right-neutral take-all)
  then have set vts = set (take (length vts - 1) vts)
    using inset
  by (metis (no-types, lifting) One-nat-def Suc-neq-Zero Suc-pred Un-insert-right
add-diff-cancel-left' butlast-conv-take diff-is-0-eq' insert-absorb len-gteq length-butlast
length-greater-0-conv list.size(3) nth-mem nth-take numeral-3-eq-3 plus-1-eq-Suc
sup-bot.right-neutral)
  then have same-set: set vts = set ?ell1
    by auto
  then have rotate i (take (length vts - 1) vts) ! 0  $\in$  set vts
    using len-gteq
  by (metis card-gteq card-length le-zero-eq length-greater-0-conv list.size(3) nth-mem
numeral-3-eq-3 zero-less-Suc)
  then have set vts = set (?ell1 @ [?ell1 ! 0])
    using same-set by auto
  then show ?thesis
    unfolding rotate-polygon-vertices-def
    using card-gteq
    by (metis butlast-conv-take)
qed

```

```

lemma arb-rotation-as-single-rotation:
  fixes  $i :: \text{nat}$ 
  shows  $\text{rotate-polygon-vertices } vts \ (\text{Suc } i) = \text{rotate-polygon-vertices } (\text{rotate-polygon-vertices } vts \ i) \ 1$ 
  unfolding rotate-polygon-vertices-def
  by (metis butlast-snoc plus-1-eq-Suc rotate-rotate)

lemma rotation-sum:
  fixes  $i \ j :: \text{nat}$ 
  shows  $\text{rotate-polygon-vertices } vts \ (i + j) = \text{rotate-polygon-vertices } (\text{rotate-polygon-vertices } vts \ i) \ j$ 
  proof (induct j)
    case 0
    thus ?case by (metis Nat.add-0-right butlast-snoc id-apply rotate0 rotate-polygon-vertices-def)
  next
    case (Suc j)
    have  $\text{rotate-polygon-vertices } vts \ (i + (\text{Suc } j)) = \text{rotate-polygon-vertices } vts \ (\text{Suc } (i + j))$  by simp
    also have ... =  $\text{rotate-polygon-vertices } (\text{rotate-polygon-vertices } vts \ (i + j)) \ 1$ 
      using arb-rotation-as-single-rotation by blast
    also have ... =  $\text{rotate-polygon-vertices } (\text{rotate-polygon-vertices } (\text{rotate-polygon-vertices } vts \ i) \ j) \ 1$ 
      using Suc.hyps by simp
    also have ... =  $\text{rotate-polygon-vertices } (\text{rotate-polygon-vertices } vts \ i) \ (\text{Suc } j)$ 
      using arb-rotation-as-single-rotation by metis
    finally show ?case .
  qed

lemma rotated-polygon-vertices-helper:
  fixes  $p :: R\text{-to-}R^2$ 
  assumes poly-p:  $\text{polygon } p$ 
  assumes p-is-path:  $p = \text{make-polygonal-path } vts$ 
  assumes p'-is:  $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$ 
  shows  $(vts ! 0) = (\text{rotate-polygon-vertices } vts \ 1) ! (\text{length } (\text{rotate-polygon-vertices } vts \ 1) - 2)$ 
     $(\text{rotate-polygon-vertices } vts \ 1) ! (\text{length } (\text{rotate-polygon-vertices } vts \ 1) - 1)$ 
     $= (vts ! 1)$ 
  proof –
    have len-gteq:  $\text{length } vts \geq 3$ 
      using polygon-at-least-3-vertices assms
      using card-length order-trans by blast
    let ?rotated-vts =  $\text{rotate-polygon-vertices } vts \ 1$ 
    have same-len:  $\text{length } ?rotated\text{-}vts = \text{length } vts$ 
      unfolding rotate-polygon-vertices-def using length-rotate
      by (metis One-nat-def Suc-pred card.empty length-append-singleton length-butlast length-greater-0-conv list.set(1) not-numeral-le-zero p-is-path poly-p polygon-at-least-3-vertices)
    then have len-rotated-gt-eq3:  $\text{length } ?rotated\text{-}vts \geq 3$ 
      using len-gteq by auto

```

```

show vts1: vts ! 0 = ?rotated-vts ! (length ?rotated-vts - 2)
  unfolding rotate-polygon-vertices-def
  using nth-rotate[of length ?rotated-vts - 2 butlast vts 1]
  Suc-diff-Suc butlast-snoc length-butlast length-greater-0-conv lessI less-nat-zero-code
  list.size(3) mod-self nat-1-add-1 nth-butlast plus-1-eq-Suc rotate-polygon-vertices-def
  same-len zero-less-diff
  by (smt (z3) One-nat-def len-gteq length-append-singleton numeral-le-one-iff
  semiring-norm(70))
  have (rotate 1 (butlast vts)) ! 0 = vts ! 1
  unfolding rotate-polygon-vertices-def
  using nth-rotate[of 0 butlast vts 1] len-gteq len-rotated-gt-eq3
  by (metis (no-types, lifting) One-nat-def Suc-le-eq length-butlast less-diff-conv
  less-nat-zero-code mod-less not-gr-zero nth-butlast numeral-3-eq-3 plus-1-eq-Suc)
  then show vts2: ?rotated-vts ! (length ?rotated-vts - 1) = vts ! 1
  unfolding rotate-polygon-vertices-def
  by (smt (verit, best) Suc-diff-Suc Suc-eq-plus1 butlast-snoc length-butlast length-greater-0-conv
  less-nat-zero-code list.size(3) nth-append-length one-add-one rotate-polygon-vertices-def
  zero-less-diff)
qed

```

lemma rotate-polygon-vertices-same-length:

```

fixes vts :: (real^2) list
assumes length vts ≥ 1
shows length vts = length (rotate-polygon-vertices vts i)
using assms
proof (induction length vts arbitrary: i)
  case 0
  then show ?case by auto
next
  case (Suc x)
  then show ?case using arb-rotation-as-single-rotation[of vts x]
  by (metis diff-Suc-1 length-append-singleton length-butlast length-rotate ro-
  tate-polygon-vertices-def)
qed

```

lemma rotated-polygon-vertices-helper2:

```

assumes len-gteq: length vts ≥ 2
assumes i < length vts - 1
assumes hd vts = last vts
shows (rotate-polygon-vertices vts 1) ! i = vts ! (i+1)
proof -
  let ?rotated-vts = rotate-polygon-vertices vts 1
  have length (butlast vts) = length vts - 1
  by auto
  then have same-len: length ?rotated-vts = length vts
  unfolding rotate-polygon-vertices-def using length-rotate len-gteq
  by (metis dual-order.trans le-add-diff-inverse length-append-singleton one-le-numeral
  plus-1-eq-Suc)
  then have len-rotated-gt-eq3: length ?rotated-vts ≥ 2

```

```

    using len-gteq by auto
  let ?n = length vts
  {assume *: i < length vts - 2
  then have same-mod: (1 + i) mod length (butlast vts) = 1+i
    using assms by simp
  have i < length (butlast vts)
    using assms by simp
  then have rotate 1 (butlast vts) ! i = butlast vts ! (i + 1)
  using nth-rotate[of i butlast vts 1] same-mod
  by (metis add.commute)
  then have (rotate-polygon-vertices vts 1) ! i = vts ! (i+1)
    by (metis (no-types, lifting) Suc-eq-plus1 <i < length (butlast vts)> butlast-snoc
length-butlast length-greater-0-conv less-nat-zero-code list.size(3) mod-less-divisor
nth-butlast plus-1-eq-Suc rotate-polygon-vertices-def same-len same-mod)
  } moreover {assume *: i = length vts - 2
  then have same-mod: (1 + i) mod length (butlast vts) = 0
    using assms
  by (metis Suc-diff-Suc <length (butlast vts) = length vts - 1> length-greater-0-conv
less-nat-zero-code list.size(3) mod-Suc mod-if one-add-one plus-1-eq-Suc zero-less-diff)
  have i < length (butlast vts)
    using assms by simp
  then have rotate-prop: rotate 1 (butlast vts) ! i = butlast vts ! 0
  using nth-rotate[of i butlast vts 1] same-mod
  by metis
  have butlast vts ! 0 = vts ! 0
    using assms(1)
  by (simp add: nth-butlast)
  then have butlast vts ! 0 = vts ! (length vts - 1)
    by (metis assms(3) hd-conv-nth last-conv-nth length-0-conv zero-diff)
  then have (rotate-polygon-vertices vts 1) ! i = vts ! (i+1)
    by (metis * rotate-prop Suc-diff-Suc Suc-eq-plus1 <butlast vts ! 0 = vts ! 0>
add-2-eq-Suc' le-add-diff-inverse2 len-gteq less-add-Suc2 one-add-one same-len but-
last-snoc length-butlast lessI nth-butlast rotate-polygon-vertices-def)
  }
  ultimately show ?thesis
    using assms(2) by linarith
qed

```

lemma *polygon-rotation-t-translation1*:

```

  assumes polygon-of p vts
  assumes p' = make-polygonal-path (rotate-polygon-vertices vts 1)
    (is p' = make-polygonal-path ?vts')
  assumes x' ∈ {(∑ i ∈ {1..k}. 1/(2i))..(∑ i ∈ {1..k+1}. 1/(2i))}
  assumes n = length vts
  assumes 0 ≤ k ∧ k ≤ n - 4
  assumes l = x' - (∑ i ∈ {1..k}. 1/(2i))
  assumes x = l/2 + (∑ i ∈ {1..(k + 1)}. 1/(2i))
  shows x ∈ {(∑ i ∈ {1..k+1}. 1/(2i))..(∑ i ∈ {1..k+2}. 1/(2i))}
    p' x' = p x

```

proof–
let $?f = \lambda(k::nat) (x::real). (x - (\sum i \in \{1..k\}. 1/(2^i))) * (2^{k+1})$
have $x \geq (\sum i \in \{1..k+1\}. 1/(2^i))$
proof–
have $l \geq 0$ **using** *assms(3,6)* **by** *auto*
then show *?thesis* **using** *assms(7)* **by** *linarith*
qed
moreover have $x \leq (\sum i \in \{1..k+2\}. 1/(2^i))$
proof–
have $x' \leq (\sum i \in \{1..k+1\}. 1/(2^i))$ **using** *assms(3)* **by** *presburger*
then have $l \leq (\sum i \in \{1..k+1\}. 1/(2^i)) - (\sum i \in \{1..k\}. 1/(2^i))$ **using**
assms(6) **by** *argo*
also have $\dots = (1/2^{k+1}) + (\sum i \in \{1..k\}. 1/(2^i)) - (\sum i \in \{1..k\}. 1/(2^i))$
using *sum-insert[of k+1 {1..k} $\lambda i. 1/(2^i)$]*
by (*smt (verit) Suc-eq-plus1 Suc-n-not-le-n add commute atLeastAtMost-Suc-conv atLeastAtMost-iff finite-atLeastAtMost le-add2 one-add-one*)
also have $\dots = (1/2^{k+1})$ **by** *argo*
finally have $l \leq (1/2^{k+1})$.
then have $x \leq (1/2^{k+1})/2 + (\sum i \in \{1..k+1\}. 1/(2^i))$ **using** *assms(7)*
by *simp*
also have $\dots = 1/2^{k+2} + (\sum i \in \{1..k+1\}. 1/(2^i))$ **by** *simp*
also have $\dots = (\sum i \in \{1..k+2\}. 1/(2^i))$
using *sum-insert[of k+2 {1..k+2} $\lambda i. 1/(2^i)$]* **by** *simp*
finally show *?thesis* .
qed
ultimately show $x: x \in \{(\sum i \in \{1..k+1\}. 1/(2^i))..(\sum i \in \{1..k+2\}. 1/(2^i))\}$
by *presburger*
have $1: n \geq 4$ **using** *polygon-vertices-length-at-least-4 assms*
using *polygon-of-def* **by** *blast*
then have $2: \text{length } vts = \text{length } ?vts'$
using *assms rotate-polygon-vertices-same-length* **by** *auto*
then have $3: \text{length } ?vts' = n$ **using** *assms* **by** *auto*

have $p' x' = ((\text{linepath } (?vts' ! k) (?vts' ! (k+1)) (?f k x'))$
using *polygon-linepath-images2[of n k ?vts' p' ?f x']* *assms(2,3,5)* $1\ 3$ **by**
fastforce
moreover have $p x = ((\text{linepath } (vts ! (k+1)) (vts ! (k+2)) (?f (k+1) x))$
using *polygon-linepath-images2[of n k+1 vts p ?f x]* *assms(2,3,5)* $1\ 2\ 3\ x$
by (*smt (verit, ccfv-threshold) Nat.diff-add-assoc add commute add-diff-cancel-left add-le-imp-le-left add-left-mono assms(1) nat-add-1-add-1 one-plus-numeral polygon-of-def semiring-norm(2) semiring-norm(4) trans-le-add1*)
moreover have $?vts' ! k = vts ! (k+1)$
using *rotated-polygon-vertices-helper2*
by (*smt (verit, best) 1 Nat.le-diff-conv2 Suc-pred' add-leD1 assms(1) assms(4) assms(5) diff-diff-cancel diff-less have-wraparound-vertex hd-conv-nth leD length-greater-0-conv less-Suc-eq nat-less-le numeral-Bit0 numeral-eq-one-iff polygon-of-def semiring-norm(83) snoc-eq-iff-butlast zero-less-numeral*)
moreover have $?vts' ! (k+1) = vts ! (k+2)$

using *rotated-polygon-vertices-helper2*[of *vts* $k+1$]
by (*metis* (*no-types*, *lifting*) *assms*(1,4,5) 1 *One-nat-def Suc-diff-Suc add-Suc-right*
diff-zero have-wraparound-vertex hd-conv-nth le-add-diff-inverse2 less-add-Suc2 nat-less-le
not-less-eq-eq numeral-Bit0 one-add-one plus-1-eq-Suc polygon-of-def snoc-eq-iff-butlast)
moreover have $?f\ k\ x' = ?f\ (k+1)\ x$ **using** *assms*(6) *assms*(7) **by** *force*
ultimately show $p'\ x' = p\ x$ **by** *presburger*
qed

lemma *polygon-rotation-t-translation1-strict*:

assumes *polygon-of* $p\ vts$
assumes $p' = \text{make-polygonal-path}\ (\text{rotate-polygon-vertices}\ vts\ 1)$
(is $p' = \text{make-polygonal-path}\ ?vts')$
assumes $x' \in \{(\sum i \in \{1..k\}. 1/(2^{\wedge}i))..<(\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))\}$
assumes $n = \text{length}\ vts$
assumes $0 \leq k \wedge k \leq n - 4$
assumes $l = x' - (\sum i \in \{1..k\}. 1/(2^{\wedge}i))$
assumes $x = l/2 + (\sum i \in \{1..(k+1)\}. 1/(2^{\wedge}i))$
shows $x \in \{(\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))..<(\sum i \in \{1..k+2\}. 1/(2^{\wedge}i))\}$
 $p'\ x' = p\ x$
proof –
let $?f = \lambda(k::nat)\ (x::real).\ (x - (\sum i \in \{1..k\}. 1/(2^{\wedge}i))) * (2^{\wedge}(k+1))$
have $x \geq (\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))$
proof –
have $l \geq 0$ **using** *assms*(3,6) **by** *auto*
then show $?thesis$ **using** *assms*(7) **by** *linarith*
qed
moreover have $x < (\sum i \in \{1..k+2\}. 1/(2^{\wedge}i))$
proof –
have $x' < (\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))$ **using** *assms*(3) **by** *auto*
then have $l < (\sum i \in \{1..k+1\}. 1/(2^{\wedge}i)) - (\sum i \in \{1..k\}. 1/(2^{\wedge}i))$ **using**
assms(6) **by** *argo*
also have $\dots = (1/2^{\wedge}(k+1)) + (\sum i \in \{1..k\}. 1/(2^{\wedge}i)) - (\sum i \in \{1..k\}. 1/(2^{\wedge}i))$
using *sum-insert*[of $k+1\ \{1..k\}\ \lambda i.\ 1/(2^{\wedge}i)$]
by (*smt* (*verit*) *Suc-eq-plus1 Suc-n-not-le-n add commute atLeastAtMost-Suc-conv atLeastAtMost-iff finite-atLeastAtMost le-add2 one-add-one*)
also have $\dots = (1/2^{\wedge}(k+1))$ **by** *argo*
finally have $l < (1/2^{\wedge}(k+1))$.
then have $x < (1/2^{\wedge}(k+1))/2 + (\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))$ **using** *assms*(7)
by *simp*
also have $\dots = 1/2^{\wedge}(k+2) + (\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))$ **by** *simp*
also have $\dots = (\sum i \in \{1..k+2\}. 1/(2^{\wedge}i))$
using *sum-insert*[of $k+2\ \{1..k+2\}\ \lambda i.\ 1/(2^{\wedge}i)$] **by** *simp*
finally show $?thesis$.
qed
ultimately show $x \in \{(\sum i \in \{1..k+1\}. 1/(2^{\wedge}i))..<(\sum i \in \{1..k+2\}. 1/(2^{\wedge}i))\}$
by *auto*
show $p'\ x' = p\ x$
using *assms*(3) *polygon-rotation-t-translation1*[OF *assms*(1) *assms*(2) - *assms*(4)]

assms(5) assms(6) assms(7)]
by (*meson atLeastAtMost-iff atLeastLessThan-iff less-eq-real-def*)
qed

lemma *polygon-rotation-t-translation2:*

assumes *polygon-of p vts*
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
(is p' = make-polygonal-path ?vts')
assumes $n = \text{length } vts$
assumes $x' \in \{(\sum i \in \{1..(n-3)\}. 1/(2^i))..(\sum i \in \{1..(n-2)\}. 1/(2^i))\}$
assumes $x = x' + 1/(2^{(n-2)})$
shows $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$
 $p' \ x' = p \ x$

proof –

let $?k = n-3$
let $?f' = (\lambda(k::nat) \ x::real. (x - (\sum i \in \{1..k\}. 1/(2^i))) * (2^{(k+1)}))$
have $n\text{-geq-4}: n \geq 4$ **using** *polygon-vertices-length-at-least-4 assms*
using *polygon-of-def by blast*
moreover then have *same-len: length vts = length ?vts'*
using *assms rotate-polygon-vertices-same-length[of vts] by auto*
moreover then have *length ?vts' = n using assms(3) by auto*
ultimately have $p'x': p' \ x' = ((\text{linepath } (?vts' \ ! \ ?k) \ (?vts' \ ! \ (?k+1)) \ (?f' \ ?k \ x')))$
using *polygon-linepath-images2[of n ?k ?vts' p' ?f' x'] assms*
by (*smt (verit, ccfv-threshold) One-nat-def Suc-diff-Suc diff-diff-left diff-is-0-eq' le-add2 le-add-diff-inverse2 linorder-not-le nat-le-linear numeral-3-eq-3 numeral-Bit0 numeral-le-iff numeral-le-one-iff numerals(1) one-plus-numeral plus-1-eq-Suc trans-le-add2*)
let $?f = (\lambda(x::real. (x - (\sum i \in \{1..n-2\}. 1/(2^i))) * (2^{(n-2)}))$
have *sum-prop: $\bigwedge i::nat. \bigwedge f::nat \Rightarrow real. (\sum i = 1..i. f \ i) + f \ (i + 1) = (\sum i = 1..i+1. f \ i)$*
by *auto*
have *sum-upto: $(\sum i = 1..n - 3. 1 / (2^i::real)) + 1 / 2^{(n-2)} = (\sum i = 1..n - 2. 1 / (2^i::real))$*
using *sum-prop[of $\lambda i. 1 / (2^i::real) \ n-3$] n-geq-4*
by (*smt (verit, del-insts) Nat.add-diff-assoc2 add-numeral-left diff-cancel2 le-add-diff-inverse le-numeral-extra(4) nat-1-add-1 nat-add-left-cancel-le numeral-Bit1 numerals(1) semiring-norm(2) semiring-norm(8) trans-le-add1*)
have $x' \geq (\sum i = 1..?k. 1 / 2^i)$
using *assms by presburger*
then have $x\text{-geq}: x \geq (\sum i \in \{1..n-2\}. 1/(2^i))$
using *assms(5) sum-upto*
by *linarith*
have $x' \leq (\sum i = 1..n - 2. 1 / 2^i)$
using *assms(4) by auto*
then have $x\text{-leq}: x \leq 1$
using *assms(5)*
by (*smt (verit, del-insts) add.left-commute add-diff-cancel-left' diff-diff-eq le-add-diff-inverse2 le-numeral-extra(4) n-geq-4 nat-add-1-add-1 numeral-Bit0 numeral-Bit1 sum-upto summation-helper trans-le-add2*)

```

show  $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$ 
  using x-geq x-leq
  by auto
  then have px:  $p\ x = (\text{linepath } (vts\ !\ (n-2))\ (vts\ !\ (n-1)))\ (?f\ x)$ 
    using polygon-linepath-images3[of n vts p x ?f] n-geq-4 assms polygon-of-def
by fastforce
  moreover have  $?vts'\ !\ (n - 3) = vts\ !\ (n-2)$ 
    using n-geq-4 assms(3) rotated-polygon-vertices-helper2 assms(1-3)
    unfolding polygon-of-def
    by (smt (verit) One-nat-def Suc-diff-Suc add.commute diff-is-0-eq diff-less
dual-order.trans have-wraparound-vertex hd-conv-nth le-add-diff-inverse length-greater-0-conv
linorder-not-le nat-1-add-1 not-add-less2 numeral-3-eq-3 plus-1-eq-Suc pos2 rotated-polygon-vertices-helper(1)
same-len snoc-eq-iff-butlast)
  moreover have  $?vts'\ !\ (n - 2) = vts\ !\ (n - 1)$ 
    using n-geq-4 assms(3) assms
    unfolding polygon-of-def
    by (metis closed-path-def list.size(3) not-numeral-le-zero polygon-def polygon-pathfinish
polygon-pathstart rotated-polygon-vertices-helper(1) same-len)
  moreover have  $?f'\ ?k\ x' = ?f\ x$  using assms(4-5) n-geq-4
    by (smt (verit, del-insts) One-nat-def Suc-diff-Suc Suc-eq-plus1 add-diff-cancel-right'
add-numeral-left le-antisym linorder-not-le numeral-3-eq-3 numeral-code(2) numeral-als(1)
semiring-norm(2) sum-upto trans-le-add2)
  ultimately show  $p'\ x' = p\ x$  using px p'x'
    by (smt (verit, ccfv-SIG) Nat.add-diff-assoc2 assms(5) diff-cancel2 le-add-diff-inverse
le-add-diff-inverse2 le-numeral-extra(4) n-geq-4 nat-1-add-1 numeral-Bit0 numeral-Bit1
trans-le-add1)
qed

```

lemma *polygon-rotation-t-translation2-strict*:

```

assumes polygon-of p vts
assumes  $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts\ 1)$ 
  ( $\text{is } p' = \text{make-polygonal-path } ?vts'$ )
assumes  $n = \text{length } vts$ 
assumes  $x' \in \{(\sum i \in \{1..(n-3)\}. 1/(2^i))..<(\sum i \in \{1..(n-2)\}. 1/(2^i))\}$ 
assumes  $x = x' + 1/(2^{(n-2)})$ 
shows  $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..<1\}$ 
   $p'\ x' = p\ x$ 
proof -
have n-geq-4:  $n \geq 4$  using polygon-vertices-length-at-least-4 assms
  using polygon-of-def by blast
have sum-prop:  $\bigwedge i::\text{nat}. \bigwedge f::\text{nat} \Rightarrow \text{real}. (\sum i = 1..i. f\ i) + f\ (i + 1) = (\sum i = 1..i+1. f\ i)$ 
  by auto
have sum-upto:  $(\sum i = 1..n - 3. 1 / (2^{\wedge} i::\text{real})) + 1 / 2^{\wedge} (n - 2) = (\sum i = 1..n - 2. 1 / (2^{\wedge} i::\text{real}))$ 
  using sum-prop[of  $\lambda i. 1 / (2^{\wedge} i::\text{real})\ n-3$ ] n-geq-4
  by (smt (verit, del-insts) Nat.add-diff-assoc2 add-numeral-left diff-cancel2 le-add-diff-inverse
le-numeral-extra(4) nat-1-add-1 nat-add-left-cancel-le numeral-Bit1 numerals(1) semir-

```

ing-norm(2) semiring-norm(8) trans-le-add1
have $x\text{-geq}: x \geq (\sum i \in \{1..n-2\}. 1/(2^i))$
using *assms(4) polygon-rotation-t-translation2[OF assms(1) assms(2) assms(3)*
- assms(5)]
by *simp*
have $x' < (\sum i = 1..n - 2. 1 / 2^i)$
using *assms(4) by auto*
then have $x\text{-leq}: x < 1$
using *assms(5)*
by (*smt (verit, del-insts) add.left-commute add-diff-cancel-left' diff-diff-eq le-add-diff-inverse2*
le-numeral-extra(4) n-geq-4 nat-add-1-add-1 numeral-Bit0 numeral-Bit1 sum-upto
summation-helper trans-le-add2)
show $x \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..<1\}$
using $x\text{-geq} x\text{-leq}$ **by** *auto*
show $p' x' = p x$
using *assms(4) polygon-rotation-t-translation2[OF assms(1) assms(2) assms(3)*
- assms(5)]
by (*meson atLeastAtMost-iff atLeastLessThan-iff less-eq-real-def*)
qed

lemma *polygon-rotation-t-translation3:*

assumes *polygon-of p vts*
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
(is p' = make-polygonal-path ?vts')
assumes $x' \in \{(\sum i \in \{1..n-2\}. 1/(2^i))..1\}$
assumes $n = \text{length } vts$
assumes $l = x' - (\sum i \in \{1..n-2\}. 1/(2^i))$
assumes $x = l * (2^{(n-3)})$
shows $x \in \{0..1/2\}$
 $p' x' = p x$

proof –

let $?f = (\lambda x::\text{real}. (x - (\sum i \in \{1..n-2\}. 1/(2^i))) * (2^{(n-2)}))$
have $n\text{-geq-4}: n \geq 4$ **using** *polygon-vertices-length-at-least-4 assms*
using *polygon-of-def by blast*
moreover then have $\text{same-len}: \text{length } vts = \text{length } ?vts'$
using *assms rotate-polygon-vertices-same-length by auto*
moreover have $\text{length-vts}': \text{length } ?vts' = n$
using *assms(4) same-len by auto*
ultimately have $p'x': p' x' = (\text{linepath } (?vts' ! (n-2)) (?vts' ! (n-1))) (?f x')$
using *polygon-linepath-images3[of n ?vts' p' x' ?f] assms*
unfolding *polygon-of-def by fastforce*

have $x\text{-is}: x = (x' - (\sum i = 1..n - 2. 1 / 2^i)) * 2^{(n-3)}$

using *assms(5-6) by auto*

then have $x\text{-gt}: x \geq 0$

using *assms(3) by simp*

have $\text{sum-prop}: k \geq 1 \implies 1 - (\sum i = 1..k. 1 / (2^i::\text{real})) = 1/(2^k)$ **for** k

proof (*induct k*)

case 0

```

    then show ?case by auto
next
case (Suc k)
{ assume *: Suc k = 1
  then have ?case by auto
} moreover
{ assume *: Suc k > 1
  then have 1 - (∑ i = 1..k. 1 / (2 ^ i::real)) = 1 / 2 ^ k
    using Suc by linarith
  then have ?case by simp
}
ultimately show ?case
  by linarith
qed
have x' ≤ 1
  using assms(3) by auto
then have x ≤ (1 - (∑ i = 1..n - 2. 1 / (2 ^ i::real))) * 2 ^ (n - 3)
  using x-is
  using mult-right-mono zero-le-power by fastforce
then have x ≤ 1 / (2 ^ (n - 2)) * 2 ^ (n - 3)
  using sum-prop n-geq-4
  by auto
then have x-lt: x ≤ 1 / 2
  using n-geq-4
  by (smt (verit, cfv-SIG) One-nat-def Suc-1 Suc-diff-Suc add-diff-cancel-right'
diff-is-0-eq dual-order.trans linorder-not-le nonzero-mult-divide-mult-cancel-right2
numeral-3-eq-3 numeral-code(2) power.simps(2) power-commutes power-not-zero
times-divide-eq-left zero-neq-numeral)
  then show x ∈ {0..1/2}
    using x-gt x-lt by auto
  moreover have n ≥ 3 using n-geq-4 by auto
  ultimately have px: p x = (linepath (vts ! 0) (vts ! 1)) (2 * x)
    using polygon-linepath-images1[of n vts] assms unfolding polygon-of-def by
blast

  have ?vts' ! (n - 2) = vts ! 0 ∧ ?vts' ! (n - 1) = vts ! 1
    unfolding rotate-polygon-vertices-def
    by (metis length-vts' assms(1) polygon-of-def rotate-polygon-vertices-def ro-
tated-polygon-vertices-helper(1) rotated-polygon-vertices-helper(2))
  moreover have ?f x' = 2 * x
  proof -
    have 2 * x = 2 * (x' - (∑ i ∈ {1..n-2}. 1 / (2 ^ i))) * (2 ^ (n - 3)) using assms
  by auto
    moreover have ... = (x' - (∑ i ∈ {1..n-2}. 1 / (2 ^ i))) * (2 ^ (n - 2))
    using n-geq-4 Suc-1 Suc-diff-Suc Suc-le-eq bot-nat-0.not-eq-extremum diff-Suc-1
le-antisym mult.left-commute mult.right-neutral mult-cancel-left not-less-eq-eq num-double
numeral-3-eq-3 numeral-eq-Suc numeral-times-numeral power.simps(2) pred-numeral-simps(2)
zero-less-diff zero-neq-numeral
  proof -

```

```

have f1:  $\forall r \text{ ra. } (ra::\text{real}) * r = r * ra$ 
  by simp
have f2:  $\forall r n \text{ ra. } (ra::\text{real}) * (r \wedge n * ra) = r \wedge \text{Suc } n * ra$ 
  by simp
have f3:  $\text{pred-numeral } (\text{num.Bit1 } \text{num.One}) = \text{Suc } (\text{Suc } 0)$ 
  by simp
have f4:  $\text{Suc } 0 = 1$ 
  by linarith
have  $\text{Suc } 1 < n$ 
  using n-geq-4 by linarith
then have  $2 * ((x' - (\sum n = 1..n - \text{Suc } 1. 1 / 2 \wedge n)) * 2 \wedge (n - 3)) =$ 
 $(x' - (\sum n = 1..n - \text{Suc } 1. 1 / 2 \wedge n)) * 2 \wedge (n - \text{Suc } 1)$ 
  using f4 f3 f2 f1 Suc-diff-Suc numeral-eq-Suc by presburger
then show ?thesis
  by (metis (no-types) Suc-1 mult.assoc)
qed
moreover have ... = ?f x' by auto
ultimately show ?thesis by presburger
qed
ultimately show  $p' x' = p x$  using p'x' px by auto
qed

```

lemma *polygon-rotation-t-translation3-strict*:

```

assumes polygon-of p vts
assumes p' = make-polygonal-path (rotate-polygon-vertices vts 1)
  (is p' = make-polygonal-path ?vts')
assumes  $x' \in \{(\sum i \in \{1..n-2\}. 1/(2 \wedge i))..<1\}$ 
assumes  $n = \text{length } vts$ 
assumes  $l = x' - (\sum i \in \{1..n-2\}. 1/(2 \wedge i))$ 
assumes  $x = l * (2 \wedge (n-3))$ 
shows  $x \in \{0..<1/2\}$ 
  p' x' = p x
proof -
  have n-geq-4:  $n \geq 4$  using polygon-vertices-length-at-least-4 assms
  using polygon-of-def by blast
  have x-is:  $x = (x' - (\sum i = 1..n - 2. 1 / 2 \wedge i)) * 2 \wedge (n - 3)$ 
  using assms(5-6) by auto
  then have x-gt:  $x \geq 0$ 
  using assms(3) by simp
  have sum-prop:  $k \geq 1 \implies 1 - (\sum i = 1..k. 1 / (2 \wedge i::\text{real})) = 1/(2 \wedge k)$  for k
  proof (induct k)
    case 0
    then show ?case by auto
  next
  case (Suc k)
  { assume *:  $\text{Suc } k = 1$ 
    then have ?case by auto
  } moreover
  { assume *:  $\text{Suc } k > 1$ 

```

```

    then have  $1 - (\sum i = 1..k. 1 / (2 ^ i::real)) = 1 / 2 ^ k$ 
      using Suc by linarith
    then have ?case by simp
  }
  ultimately show ?case
    by linarith
qed
have  $x' < 1$ 
  using assms(3) by auto
then have  $x < (1 - (\sum i = 1..n - 2. 1 / (2 ^ i::real))) * 2 ^ (n - 3)$ 
  using x-is
  using mult-right-mono zero-le-power by fastforce
then have  $x < 1/(2^{n-2})*2^{n-3}$ 
  using sum-prop n-geq-4
  by auto
then have x-lt:  $x < 1/2$ 
  using n-geq-4
  by (smt (verit, ccfv-SIG) One-nat-def Suc-1 Suc-diff-Suc add-diff-cancel-right'
diff-is-0-eq dual-order.trans linorder-not-le nonzero-mult-divide-mult-cancel-right2
numeral-3-eq-3 numeral-code(2) power.simps(2) power-commutes power-not-zero
times-divide-eq-left zero-neq-numeral)
  show  $x \in \{0..<1/2\}$ 
    using x-lt x-gt by auto
  show  $p' x' = p x$ 
    using assms(3) polygon-rotation-t-translation3[OF assms(1) assms(2) - assms(4)
assms(5) assms(6)]
    by simp
qed

lemma f-gteq-0-sum-gt:  $\bigwedge f::nat \Rightarrow real. (\bigwedge i::nat. (f i) > 0) \Longrightarrow a > b \Longrightarrow (\sum i = 1..a. (f i)) > (\sum i = 1..b. (f i))$  for  $a b :: nat$ 
proof (induct a arbitrary: b)
  case 0
  then show ?case by auto
next
  case (Suc a)
  {assume *:  $b = a$ 
  then have  $sum f \{1..(Suc a)\} = sum f \{1.. b\} + f (Suc a)$ 
    by force
  then have ?case
    using Suc(2)[of Suc a] * by linarith
  } moreover {assume *:  $b < a$ 
  then have ?case using Suc
  by (smt (verit, ccfv-threshold) Suc-eq-plus1 dual-order.trans le-add2 sum.nat-ivl-Suc')}
  }
  ultimately show ?case
    using Suc.prem(2) less-antisym by blast
qed

```

lemma *rotation-intervals-disjoint*:

assumes $k1 \neq k2$
shows $\{\sum i = 1..k1. 1 / (2^{\wedge} i::real)..<\sum i = 1..k1+1. 1 / 2^{\wedge} i\} \cap \{\sum i = 1..k2. 1 / (2^{\wedge} i::real)..<\sum i = 1..k2+1. 1 / 2^{\wedge} i\} = \{\}$
proof –
have *lambda-gt*: $(\wedge i. 0 < 1 / (2^{\wedge} i::real))$
by *simp*
have *h1*: *?thesis* **if** $*:k1 < k2$
proof –
have *eo*: $k1+1 \leq k2$
using $*$ **by** *auto*
have $k1+1 = k2 \implies (\sum i = 1..k1+1. 1 / 2^{\wedge} i) \leq (\sum i = 1..k2. 1 / (2^{\wedge} i::real))$
by *auto*
have $(\sum i = 1..k1+1. 1 / 2^{\wedge} i) \leq (\sum i = 1..k2. 1 / (2^{\wedge} i::real))$ **if** $**:$
 $k1+1 < k2$
using *f-gteq-0-sum-gt*[*OF lambda-gt ***]
using *less-eq-real-def* **by** *presburger*
then **have** $(\sum i = 1..k1+1. 1 / 2^{\wedge} i) \leq (\sum i = 1..k2. 1 / (2^{\wedge} i::real))$
using $*$ *eo* **by** *fastforce*
then **show** *?thesis* **by** *auto*
qed
have *h2*: *?thesis* **if** $*: k2 < k1$
proof –
have *eo*: $k2+1 \leq k1$
using $*$ **by** *auto*
have $k2+1 = k1 \implies (\sum i = 1..k2+1. 1 / 2^{\wedge} i) \leq (\sum i = 1..k1. 1 / (2^{\wedge} i::real))$
by *auto*
have $(\sum i = 1..k2+1. 1 / 2^{\wedge} i) \leq (\sum i = 1..k1. 1 / (2^{\wedge} i::real))$ **if** $**:$
 $k2+1 < k1$
using *f-gteq-0-sum-gt*[*OF lambda-gt ***]
using *less-eq-real-def* **by** *presburger*
then **have** $(\sum i = 1..k2+1. 1 / 2^{\wedge} i) \leq (\sum i = 1..k1. 1 / (2^{\wedge} i::real))$
using $*$ *eo* **by** *fastforce*
then **show** *?thesis* **by** *auto*
qed
show *?thesis*
using *h1 h2 assms* **by** *linarith*
qed

lemma *bounding-interval-helper1*:

shows $(\sum i = 1..k. 1 / (2^{\wedge} i::real)) = (2^{\wedge} k - 1) / (2^{\wedge} k)$
proof(*induct k*)
case 0
then **show** *?case* **by** *simp*
next
case (*Suc k*)
have $(\sum i = 1..(Suc k). 1 / (2^{\wedge} i::real)) = (\sum i = 1..k. 1 / (2^{\wedge} i::real)) +$

$1/2^{\wedge}(Suc\ k)$
by force
also have ... = $(2^{\wedge}k - 1)/(2^{\wedge}k) + 1/2^{\wedge}(Suc\ k)$ **using** *Suc.hyps* **by** *presburger*
also have ... = $(2^{\wedge}k - 1)/(2^{\wedge}k) + 1/2^{\wedge}(k+1)$ **by** *simp*
also have ... = $(2^{\wedge}(k+1) - 1)/(2^{\wedge}(k+1))$
by (*smt (verit, del-insts) Suc add.commute add-diff-cancel-right' add-divide-distrib calculation field-sum-of-halves le-add2 plus-1-eq-Suc power-divide power-one summation-helper*)
finally show *?case* **by force**
qed

lemma *bounding-interval-helper2*:

fixes $x :: real$
assumes $x \in \{0..<1\}$
shows $\exists k. x < (\sum i = 1..k. 1 / (2^{\wedge} i :: real))$
proof –
let $?f = \lambda k :: nat. (2^{\wedge}k - 1)/(2^{\wedge}k)$
have *lim*: $\forall \varepsilon :: real > 0. \exists k. (1 - (?f\ k)) < \varepsilon$
proof *clarify*
fix $\varepsilon :: real$
assume $\varepsilon > 0$
then obtain m **where** $m > 0 \wedge 1 / m < \varepsilon$
by (*metis Groups.mult-ac(2) divide-less-eq linordered-field-no-ub order-less-trans zero-less-divide-1-iff*)
moreover obtain k **where** $2^{\wedge}k > m$ **using** *real-arch-pow* **by** *fastforce*
ultimately have $1 / (2^{\wedge}k) < \varepsilon$ **by** (*smt (verit) frac-less2*)
moreover have $(1 :: real) - ((2^{\wedge}k - 1) / (2^{\wedge}k)) = (1 / (2^{\wedge}k))$ **by** (*simp add: diff-divide-distrib*)
ultimately show $\exists k. 1 - (2^{\wedge}k - 1) / (2^{\wedge}k) < \varepsilon$ **by** (*smt (verit)*)
qed
have $\exists k. ?f\ k > x$
proof –
let $?e = 1 - x$
obtain k **where** $1 - (?f\ k) < ?e$ **by** (*metis assms lim atLeastLessThan-iff diff-gt-0-iff-gt*)
thus *?thesis* **by auto**
qed
thus *?thesis* **using** *bounding-interval-helper1* **by** *presburger*
qed

lemma *bounding-interval-for-reals-btw01*:

fixes $x :: real$
assumes $x \in \{0..<1\}$
shows $\exists k. x \in \{(\sum i \in \{1..k\}. 1/(2^{\wedge}i :: real))..<(\sum i \in \{1..(k+1)\}. 1/(2^{\wedge}i))\}$
proof –
let $?S = \lambda k. (\sum i = 1..k. 1 / (2^{\wedge} i :: real))$
let $?A = \{k :: nat. x < (\sum i = 1..k. 1 / (2^{\wedge} i :: real))\}$
let $?m = LEAST\ k. k \in ?A$
have $\exists k. x < (\sum i = 1..k. 1 / (2^{\wedge} i :: real))$ **using** *assms bounding-interval-helper2*

by *blast*
then have $?m \in ?A$ **by** (*metis (mono-tags, lifting) LeastI2-wellorder mem-Collect-eq*)
moreover then have $?m - 1 \notin ?A$
by (*smt (verit, ccfv-SIG) One-nat-def Suc-n-not-le-n Suc-pred' assms atLeast-LessThan-iff atLeastatMost-empty' bot-nat-0.not-eq-extremum linorder-not-less mem-Collect-eq not-less-Least sum.empty*)
ultimately have $x < (\sum i = 1..?m. 1 / (2 \wedge i::real)) \wedge x \geq (\sum i = 1..?m-1. 1 / (2 \wedge i::real))$
by *simp*
thus *?thesis*
by (*smt (verit, best) add commute assms atLeastLessThan-iff le-add-diff-inverse linorder-not-less sum.head-if*)
qed

lemma *all-rotation-intervals-between-0and1*:
shows $\{(\sum i \in \{1..k\}. 1/(2 \wedge i::real))..(\sum i \in \{1..(k+1)\}. 1/(2 \wedge i))\} \subseteq \{0..<1\}$
proof –
have *gt*: $\bigwedge k. (\sum i \in \{1..k\}. 1/(2 \wedge i::real)) \geq 0$
by (*simp add: sum-nonneg*)
have *lt*: $\bigwedge k. (\sum i \in \{1..k\}. 1/(2 \wedge i::real)) < 1$
by (*smt (verit, ccfv-SIG) diff-Suc-1 f-gteq-0-sum-gt less-Suc-eq-le linorder-not-le summation-helper zero-less-divide-1-iff zero-less-power*)
show *?thesis*
using *gt lt*
by (*meson atLeastAtMost-subseteq-atLeastLessThan-iff*)
qed

lemma *all-rotation-intervals-between-0and1-strict*:
shows $\{(\sum i \in \{1..k\}. 1/(2 \wedge i::real))..<(\sum i \in \{1..(k+1)\}. 1/(2 \wedge i))\} \subseteq \{0..<1\}$
using *all-rotation-intervals-between-0and1*
by (*smt (verit, ccfv-SIG) atLeastAtMost-subseteq-atLeastLessThan-iff ivl-subset nle-le order-trans*)

lemma *one-polygon-rotation-is-loop-free*:
assumes *polygon-of p vts*
assumes $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$
(is $p' = \text{make-polygonal-path } ?vts')$
shows *loop-free p'*
proof(*rule ccontr*)
assume $\neg \text{loop-free } p'$
moreover have $p' \ 0 = p' \ 1$
using *assms*
by (*smt (verit, ccfv-SIG) assms(2) butlast-snoc length-butlast linepath-0' linepath-1' make-polygonal-path.simps(1) not-gr-zero nth-append-length nth-butlast path-defs(2) path-defs(3) polygon-pathfinish polygon-pathstart rotate-polygon-vertices-def*)
ultimately obtain $x' \ y'$ **where** $x' \ y': x' < y' \wedge \{x', y'\} \subseteq \{0..<1\} \wedge p' \ x' = p' \ y'$
unfolding *loop-free-def*
by (*smt (verit, del-insts) atLeastAtMost-iff atLeastLessThan-iff bot-least in-*)

sert-subset linorder-not-le order.refl order-antisym zero-less-one)

```

let ?n = length vts
have n-geq-4: ?n ≥ 4 using polygon-vertices-length-at-least-4 assms
  using polygon-of-def by blast
obtain xk where x'-in: x' ∈ {∑ i ∈ {1..xk}. 1/(2i)}..<(∑ i ∈ {1..(xk + 1)}.
1/(2i))} using x'y'
  using bounding-interval-for-reals-btw01 x'y'
  by (metis insert-subset)
then have xk-gteq: xk ≥ 0
  by blast
obtain yk where y'-in: y' ∈ {∑ i ∈ {1..yk}. 1/(2i)}..<(∑ i ∈ {1..(yk + 1)}.
1/(2i))}
  using bounding-interval-for-reals-btw01 x'y'
  by (metis insert-subset)
then have yk-gteq: yk ≥ 0
  by blast

have all-pows-of-2-pos: (∧ i. 0 < 1 / (2i :: real))
  by simp

```

```

let ?x1 = (x' - (∑ i ∈ {1..xk}. 1/(2i)))/2 + (∑ i ∈ {1..(xk + 1)}. 1/(2i))
have xk-lt-nminus3: xk ≤ ?n - 4 ⇒ ?x1 ∈ {∑ i ∈ {1..xk+1}. 1/(2i)}..<(∑ i
∈ {1..xk+2}. 1/(2i))} ∧ p ?x1 = p' x'
  using polygon-rotation-t-translation1-strict[OF assms(1) assms(2) x'-in] xk-gteq
  by metis
let ?y1 = (y' - (∑ i ∈ {1..yk}. 1/(2i)))/2 + (∑ i ∈ {1..(yk + 1)}. 1/(2i))
have yk-lt-nminus3: yk ≤ ?n - 4 ⇒ ?y1 ∈ {∑ i ∈ {1..yk+1}. 1/(2i)}..<(∑ i
∈ {1..yk+2}. 1/(2i))} ∧ p ?y1 = p' y'
  using polygon-rotation-t-translation1-strict[OF assms(1) assms(2) y'-in] yk-gteq

by metis

```

```

let ?x2 = x' + 1/(2(?n-2))
have xk = ?n-3 ⇒ x' ∈ {∑ i = 1..length vts - 3. 1 / (2i :: real)}..<∑ i =
1..length vts - 2. 1 / 2i}
  using x'-in
  by (smt (verit, best) Nat.add-diff-assoc2 ‹4 ≤ length vts› diff-cancel2 le-add-diff-inverse
nat-add-left-cancel-le nat-le-linear numeral-Bit0 numeral-Bit1 numerals(1) trans-le-add1)
then have xk-eq-nminus3: xk = ?n - 3 ⇒ p ?x2 = p' x' ∧ ?x2 ∈ {∑ i ∈
{1..?n-2}. 1/(2i)}..<1}
  using polygon-rotation-t-translation2-strict[OF assms(1) assms(2), of ?n x'
?x2] x'-in xk-gteq
  by presburger
let ?y2 = y' + 1/(2(?n-2))
have yk = ?n-3 ⇒ y' ∈ {∑ i = 1..length vts - 3. 1 / (2i :: real)}..<∑ i =
1..length vts - 2. 1 / 2i}
  using y'-in

```

```

    by (smt (verit, best) Nat.add-diff-assoc2 <4 ≤ length vts> diff-cancel2 le-add-diff-inverse
    nat-add-left-cancel-le nat-le-linear numeral-Bit0 numeral-Bit1 numerals(1) trans-le-add1)
    then have yk-eq-nminus3: yk = ?n - 3 ⇒ p ?y2 = p' y' ∧ ?y2 ∈ {(∑ i ∈
    {1..?n-2}. 1/(2i))..<1}
      using polygon-rotation-t-translation2-strict[OF assms(1) assms(2), of ?n y'
    ?y2] x'-in xk-gteq
      by presburger

    let ?x3 = (x' - (∑ i ∈ {1..?n-2}. 1/(2i)))*(2?n-3)
    have x'-leq: x' < 1
      using x'y' by simp
    have x'-geq: xk ≥ ?n - 2 ⇒ (∑ i = 1..xk. 1 / (2i)) ≥ (∑ i = 1..length
    vts - 2. 1 / (2i))
      using x'-in f-gteq-0-sum-gt[of λi. 1 / (2i)]
      by (metis le-antisym less-eq-real-def linorder-not-le zero-less-divide-1-iff zero-less-numeral
    zero-less-power)
    have xk ≥ ?n-2 ⇒ x' ∈ {∑ i = 1..length vts - 2. 1 / (2i)}
      using x'-leq x'-geq x'-in
      by fastforce
    then have xk-gt-nminus3: xk ≥ ?n - 2 ⇒ p ?x3 = p' x' ∧ ?x3 ∈ {0..<1/2}
      using polygon-rotation-t-translation3-strict[OF assms(1) assms(2), of x' ?n]
    xk-gteq
      by presburger
    let ?y3 = (y' - (∑ i ∈ {1..?n-2}. 1/(2i)))*(2?n-3)
    have y'-leq: y' < 1
      using x'y' by simp
    have y'-geq: yk ≥ ?n - 2 ⇒ (∑ i = 1..yk. 1 / (2i)) ≥ (∑ i = 1..length
    vts - 2. 1 / (2i))
      using y'-in f-gteq-0-sum-gt[of λi. 1 / (2i)]
      by (metis le-antisym less-eq-real-def linorder-not-le zero-less-divide-1-iff zero-less-numeral
    zero-less-power)
    have yk ≥ ?n-2 ⇒ y' ∈ {∑ i = 1..length vts - 2. 1 / (2i)}
      using y'-leq y'-geq y'-in
      by fastforce
    then have yk-gt-nminus3: yk ≥ ?n - 2 ⇒ p ?y3 = p' y' ∧ ?y3 ∈ {0..<1/2}
      using polygon-rotation-t-translation3-strict[OF assms(1) assms(2), of y' ?n]
    yk-gteq
      by presburger

    have interval-helper: a1 ≥ b2 ∧ x ∈ {a1..<a2} ∧ y ∈ {b1..<b2} ⇒ y < x for
    a1 a2 b1 b2 x y::real
      by simp

    { assume xk-lt: xk < ?n - 3
      then have p-x': p ?x1 = p' x'
        using xk-lt-nminus3 by auto
      have x1-in: ?x1 ∈ {(∑ i ∈ {1..(xk + 1)}. 1/(2i))..<(∑ i ∈ {1..(xk + 2)}.
    1/(2i))}
        using xk-lt xk-lt-nminus3

```

```

    by auto
  then have x1-in-01: ?x1 ∈ {0..<1}
    using all-rotation-intervals-between-0and1-strict[of xk+1]
    by fastforce
  { assume yk-lt: yk < ?n - 3
    then have p-y': p ?y1 = p' y'
      using yk-lt-nminus3 by auto
    have y1-in: ?y1 ∈ {(∑ i ∈ {1..(yk + 1)}. 1/(2i))..<(∑ i ∈ {1..(yk + 2)}.
1/(2i))}
      using yk-lt yk-lt-nminus3 by auto
    then have y1-in-01: ?y1 ∈ {0..<1}
      using all-rotation-intervals-between-0and1-strict[of yk+1]
      by fastforce
    have {∑ i = 1..xk + 1. 1 / 2i..<∑ i = 1..xk + 2. 1 / (2i::real)} ∩ {∑ i
= 1..yk + 1. 1 / (2i::real)..<∑ i = 1..yk + 2. 1 / 2i} = {} if xk-neq:xk ≠
yk
      using rotation-intervals-disjoint[of xk+1 yk+1] xk-neq
      by fastforce
    then have eq-then-eq: ?x1 = ?y1 ⇒ xk = yk
      using x1-in y1-in
      by (smt (verit) Int-iff empty-iff)
    have xk = yk ⇒ ?x1 ≠ ?y1
      using x'y' x1-in y1-in by simp
    then have ?x1 ≠ ?y1
      using eq-then-eq by blast
    moreover have {?x1, ?y1} ⊆ {0..<1}
      using x1-in-01 y1-in-01 by fast
    ultimately have ?x1 ≠ ?y1 ∧ {?x1, ?y1} ⊆ {0..<1} ∧ p ?x1 = p ?y1
      using p-x' p-y' x'y' by presburger
    then have ∃ x y . x ≠ y ∧ {x, y} ⊆ {0..<1} ∧ p x = p y
      by auto
    then have False
      using asms(1) unfolding polygon-of-def polygon-def simple-path-def loop-free-def
      by fastforce
  } moreover { assume yk = ?n - 3
    then have y2: p ?y2 = p' y' ∧ ?y2 ∈ {(∑ i ∈ {1..?n-2}. 1/(2i))..<1}
      using yk-eq-nminus3
      by auto
    then have y2-in-01: ?y2 ∈ {0..<1}
      using all-rotation-intervals-between-0and1-strict[of ?n-2]
      by fastforce
    have xkplus-eq: xk + 2 = ?n - 2 ⇒ (∑ i ∈ {1..(xk + 2)}. 1/(2i::real)) ≤
(∑ i ∈ {1..?n-2}. 1/(2i))
      by simp
    have xkplus-lt: xk + 2 < ?n - 2 ⇒ (∑ i ∈ {1..(xk + 2)}. 1/(2i::real)) ≤
(∑ i ∈ {1..?n-2}. 1/(2i))
      using xk-lt f-gteq-0-sum-gt[OF all-pows-of-2-pos, of xk + 2 ?n - 2]
      by (smt (verit, best) f-gteq-0-sum-gt zero-less-divide-1-iff zero-less-power)
    then have (∑ i ∈ {1..(xk + 2)}. 1/(2i::real)) ≤ (∑ i ∈ {1..?n-2}. 1/(2i))

```

```

    using xkplus-eq xkplus-lt xk-lt
    using One-nat-def Suc-diff-Suc Suc-eq-plus1 Suc-le-eq add-Suc-right le-neq-implies-less
    linorder-not-le nat-1-add-1 nat-diff-split numeral-3-eq-3 xk-gteq by linarith
    then have  $?x1 \neq ?y2$ 
      using x1-in y2
      by (smt (verit, ccfv-SIG) interval-helper)
    moreover have  $\{?x1, ?y2\} \subseteq \{0..<1\}$ 
      using x1-in-01 y2-in-01 by fast
    ultimately have  $?x1 \neq ?y2 \wedge \{?x1, ?y2\} \subseteq \{0..<1\} \wedge p ?x1 = p ?y2$ 
      using p-x' y2 x'y' by presburger
    then have  $\exists x y . x \neq y \wedge \{x, y\} \subseteq \{0..<1\} \wedge p x = p y$ 
      by auto
    then have False
      using assms(1) unfolding polygon-of-def polygon-def simple-path-def
loop-free-def
      by fastforce
  }
  moreover { assume  $yk > ?n - 3$ 
    then have  $y3: p ?y3 = p' y' \wedge ?y3 \in \{0..<(1/2::real)\}$ 
      using yk-gt-nminus3
      by auto
    then have y3-in-01:  $?y3 \in \{0..<1\}$ 
      by simp

    have simplify-interval:  $(\sum i = 1..1. 1 / (2 \wedge i::real)) = 1/2$ 
      by simp
    then have xk-eq-0:  $xk = 0 \implies (\sum i \in \{1..(xk + 1)\}. 1/(2 \wedge i::real)) \geq 1/2$ 
      by simp
    have  $xk > 0 \implies (\sum i \in \{1..(xk + 1)\}. 1/(2 \wedge i::real)) \geq 1/2$ 
      using f-gteq-0-sum-gt[OF all-pows-of-2-pos, of 1 xk + 1]
      simplify-interval
      by (smt (verit, ccfv-SIG) Suc-le-eq add.commute add.right-neutral all-pows-of-2-pos
f-gteq-0-sum-gt linorder-not-le plus-1-eq-Suc)
    then have  $(\sum i \in \{1..(xk + 1)\}. 1/(2 \wedge i::real)) \geq 1/2$ 
      using xk-eq-0 xk-gteq by blast
    then have  $?x1 \neq ?y3$ 
      using x1-in y3
      by (smt (verit, best) interval-helper)
    moreover have  $\{?x1, ?y3\} \subseteq \{0..<1\}$ 
      using x1-in-01 y3-in-01 by fast
    ultimately have  $?x1 \neq ?y3 \wedge \{?x1, ?y3\} \subseteq \{0..<1\} \wedge p ?x1 = p ?y3$ 
      using p-x' y3 x'y'
      by presburger
    then have  $\exists x y . x \neq y \wedge \{x, y\} \subseteq \{0..<1\} \wedge p x = p y$ 
      by auto
    then have False
      using assms(1) unfolding polygon-of-def polygon-def simple-path-def
loop-free-def
      by fastforce
  }

```

```

}
ultimately have False by linarith
} moreover { assume xk-eq : xk = ?n-3
then have p-x': p ?x2 = p' x'
using xk-eq-nminus3 by auto
have x2-in: ?x2 ∈ {(∑ i ∈ {1..?n-2}. 1/(2i))..<1}
using xk-eq xk-eq-nminus3
by auto
then have ?x2 ≥ 0
using n-geq-4
by (metis add-sign-intros(4) atLeastLessThan-iff insert-subset leD nle-le
power-one-over x'y' zero-le-power zero-less-divide-1-iff zero-less-numeral)
then have x2-in-01: ?x2 ∈ {0..<1}
using x2-in by auto
{ assume yk < ?n - 3
then have interval-helper-helper: (∑ i = 1..yk + 1. 1 / (2i :: real)) ≤ (∑ i
= 1..xk. 1 / (2i :: real))
using xk-eq f-gteq-0-sum-gt
by (metis Suc-eq-plus1 less-eq-real-def linorder-neqE-nat not-less-eq zero-less-divide-1-iff
zero-less-numeral zero-less-power)
then have x' > y'
using x'-in y'-in interval-helper[of (∑ i = 1..yk + 1. 1 / (2i :: real))
(∑ i = 1..xk. 1 / (2i :: real))]
by blast
then have False using x'y'
by auto
} moreover { assume yk = ?n - 3
then have y2: p ?y2 = p' y' ∧ ?y2 ∈ {(∑ i ∈ {1..?n-2}. 1/(2i))..<1}
using yk-eq-nminus3
by auto
then have y2-in-01: ?y2 ∈ {0..<1}
using all-rotation-intervals-between-0and1-strict[of ?n-2]
by fastforce
then have ?x2 ≠ ?y2
using x'y' by auto
moreover have {?x2, ?y2} ⊆ {0..<1}
using x2-in-01 y2-in-01 by fast
ultimately have ?x2 ≠ ?y2 ∧ {?x2, ?y2} ⊆ {0..<1} ∧ p ?x2 = p ?y2
using p-x' y2 x'y' by presburger
then have ∃ x y . x ≠ y ∧ {x, y} ⊆ {0..<1} ∧ p x = p y
by meson
then have False
using assms(1) unfolding polygon-of-def polygon-def simple-path-def
loop-free-def
by fastforce
} moreover { assume yk-gt: yk > ?n - 3
then have y3: p ?y3 = p' y'
using yk-gt-nminus3 by auto
have y3-in: ?y3 ∈ {0..<1/2}

```

```

    using yk-gt yk-gt-nminus3
    by auto
  then have y3-in-01:  $?y3 \in \{0..<1\}$ 
    by auto
  have  $(\sum i = 1..length\ vts - 2. 1 / (2 \wedge i::real)) > (\sum i = 1..1. 1 / (2 \wedge i::real))$ 
    using n-geq-4 f-gteq-0-sum-gt[OF all-pows-of-2-pos,of 1 length vts - 2]
    by fastforce
  then have  $(\sum i = 1..length\ vts - 2. 1 / (2 \wedge i::real)) > 1/2$ 
    by simp
  then have  $?x2 \neq ?y3$ 
    using y3-in x2-in by auto
  moreover have  $\{?x2, ?y3\} \subseteq \{0..<1\}$ 
    using x2-in-01 y3-in-01 by fast
  ultimately have  $?x2 \neq ?y3 \wedge \{?x2, ?y3\} \subseteq \{0..<1\} \wedge p\ ?x2 = p\ ?y3$ 
    using p-x' y3 x'y' by presburger
  then have  $\exists x\ y. x \neq y \wedge \{x, y\} \subseteq \{0..<1\} \wedge p\ x = p\ y$ 
    by meson
  then have False
    using assms(1) unfolding polygon-of-def polygon-def simple-path-def loop-free-def
    by fastforce
}
ultimately have False
  using not-less-iff-gr-or-eq by auto
} moreover { assume xk-gt:  $xk > ?n - 3$ 
  then have  $p\ x': p\ ?x3 = p'\ x'$ 
    using xk-gt-nminus3 by auto
  have x3-in:  $?x3 \in \{0..<1/2\}$ 
    using xk-gt xk-gt-nminus3
    by auto
  then have x3-in-01:  $?x3 \in \{0..<1\}$ 
    by auto
  { assume  $yk \leq ?n - 3$ 
    then have  $(\sum i = 1..xk. 1 / (2 \wedge i::real)) \geq (\sum i = 1..yk + 1. 1 / (2 \wedge i::real))$ 
      using xk-gt f-gteq-0-sum-gt[of  $\lambda i. 1 / (2 \wedge i::real)$  xk yk]
      proof -
        obtain rr :: nat  $\Rightarrow$  real where
          f1:  $\forall B\ x. rr\ B\ x = 1 / 2 \wedge B\ x$ 
          by force
        then have f2:  $\forall n. 0 < rr\ n$ 
          by simp
        have  $yk < xk$ 
          using  $\langle length\ vts - 3 < xk \rangle \langle yk \leq length\ vts - 3 \rangle$  order-le-less-trans by
          blast
        then show thesis
          using f2 f1 by (metis (no-types) Suc-eq-plus1 f-gteq-0-sum-gt less-eq-real-def nat-neq-iff not-less-eq order.refl)
      }
  }
}

```

```

qed
then have  $x' > y'$ 
  using  $x'$ -in  $y'$ -in interval-helper[of  $(\sum i = 1..yk + 1. 1 / (2^i::real)) (\sum i = 1..xk. 1 / (2^i::real))$ ]
  by blast
then have False using  $x'y'$ 
  by auto
} moreover
{ assume  $yk$ -gt:  $yk > ?n - 3$ 
  then have  $p$ - $y'$ :  $p ?y3 = p' y'$ 
    using  $yk$ -gt- $n$ minus3 by auto
  have  $y3$ -in:  $?y3 \in \{0..<1/2\}$ 
    using  $yk$ -gt  $yk$ -gt- $n$ minus3
    by auto
  then have  $y3$ -in-01:  $?y3 \in \{0..<1\}$ 
    by auto
  have  $(x' - (\sum i = 1..length vts - 2. 1 / 2^i)) \neq$ 
     $(y' - (\sum i = 1..length vts - 2. 1 / 2^i))$ 
    using  $x'y'$  by auto
  then have  $?x3 \neq ?y3$  by auto
  moreover have  $\{?x3, ?y3\} \subseteq \{0..<1\}$ 
    using  $x3$ -in-01  $y3$ -in-01 by fast
  ultimately have  $?x3 \neq ?y3 \wedge \{?x3, ?y3\} \subseteq \{0..<1\} \wedge p ?x3 = p ?y3$ 
    using  $p$ - $x'$   $p$ - $y'$   $x'y'$ 
    by presburger
  then have  $\exists x y . x \neq y \wedge \{x, y\} \subseteq \{0..<1\} \wedge p x = p y$ 
    by meson
  then have False
    using assms(1) unfolding polygon-of-def polygon-def simple-path-def
loop-free-def
    by fastforce
}
ultimately have False by linarith
}
ultimately show False by linarith
qed

```

lemma *one-rotation-is-polygon*:

fixes $p :: R$ -to- R^2

fixes $i :: \text{nat}$

assumes $poly$ - p : *polygon* p **and**

p -is-path: $p = \text{make-polygonal-path } vts$ **and**

p' -is: $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$

(**is** $p' = \text{make-polygonal-path } ?vts'$)

shows *polygon* p'

proof –

have *polygonal-path* p' **using** p' -is **by** (*simp add: polygonal-path-def*)

moreover have *closed-path* p'

using p' -is **unfolding** *rotate-polygon-vertices-def* *closed-path-def*

by (*metis* (*no-types*, *opaque-lifting*) *Nil-is-append-conv* *append-self-conv2* *diff-Suc-1* *hd-append2* *hd-conv-nth* *length-append-singleton* *make-polygonal-path-gives-path* *not-Cons-self* *nth-Cons-0* *nth-append-length* *pathfinish-def* *pathstart-def* *polygon-pathfinish* *polygon-pathstart*)
moreover have *simple-path p'*
using *one-polygon-rotation-is-loop-free*
by (*metis* *make-polygonal-path-gives-path* *p'-is* *p-is-path* *poly-p* *polygon-of-def* *simple-path-def*)
ultimately show *?thesis unfolding polygon-def by simp*
qed

lemma *rotation-is-polygon*:
fixes *p :: R-to-R2*
fixes *i :: nat*
assumes *polygon p* **and**
 $p = \text{make-polygonal-path } vts$
shows *polygon (make-polygonal-path (rotate-polygon-vertices vts i))*
using *assms*
proof (*induct i*)
case *0*
then show *?case using rotate0 unfolding rotate-polygon-vertices-def*
by (*smt* (*z3*) *assms(2)* *butlast.simps(1)* *butlast-conv-take* *eq-id-iff* *have-wraparound-vertex* *hd-append2* *hd-conv-nth* *rotate-polygon-vertices-def* *rotate-polygon-vertices-same-set* *self-append-conv2* *the-elem-set*)
next
case (*Suc i*)
then show *?case using one-rotation-is-polygon arb-rotation-as-single-rotation*
by *metis*
qed

lemma *polygon-rotate-mod*:
fixes *vts :: (real^2) list*
assumes $n = \text{length } vts$
assumes $n \geq 2$
assumes $\text{hd } vts = \text{last } vts$
shows $\text{rotate-polygon-vertices } vts (n - 1) = vts$
proof –
let $?vts' = \text{rotate } (n - 1) (\text{butlast } vts)$
have $\text{rotate-polygon-vertices } vts (n - 1) = ?vts' @ [?vts!0]$
unfolding *rotate-polygon-vertices-def* **by** *metis*
moreover have $?vts' = \text{butlast } vts$ **using** *assms* **by** *simp*
moreover have $\dots = \text{rotate } 0 (\text{butlast } vts)$ **by** *simp*
moreover then have $\dots @ [!0] = \text{rotate-polygon-vertices } vts 0$
unfolding *rotate-polygon-vertices-def* **by** *metis*
moreover have $\dots = vts$
unfolding *rotate-polygon-vertices-def* **using** *assms*
by (*metis* (*no-types*, *lifting*) *Suc-le-eq* *calculation(3)* *hd-conv-nth* *length-butlast* *length-greater-0-conv* *nat-1-add-1* *nth-butlast* *order-less-le-trans* *plus-1-eq-Suc* *pos2* *snoc-eq-iff-butlast* *zero-less-diff*)

ultimately show *?thesis* **by** *argo*
qed

lemma *polygon-rotate-mod-arb*:

fixes *vts* :: (real²) list
assumes *n = length vts*
assumes *n ≥ 2*
assumes *hd vts = last vts*
shows *rotate-polygon-vertices vts ((n - 1) * i) = vts*
proof(*induct i*)
case 0
then show *?case* **using** *polygon-rotate-mod*
by (*metis append.right-neutral append-Nil assms(1) assms(2) assms(3) id-apply length-butlast mult-zero-right rotate0 rotate-append rotate-polygon-vertices-def*)
next
case (*Suc i*)
then have *vts = rotate-polygon-vertices vts ((n - 1) * i)* **using** *Suc.prem*s **by** *argo*
also have *... = rotate-polygon-vertices vts ((n - 1) * Suc i)*
using *polygon-rotate-mod assms(1) assms(2) assms(3) calculation rotation-sum*
by (*metis mult-Suc-right*)
finally show *?case* **by** *argo*
qed

lemma *unrotation-is-polygon*:

fixes *p* :: R-to-R²
fixes *i*:: nat
assumes *polygon (make-polygonal-path (rotate-polygon-vertices vts i))*
(is *polygon (make-polygonal-path ?vts')*
p = make-polygonal-path vts
hd vts = last vts
shows *polygon p*
proof–
have *len-vts: length vts ≥ 2*
using *assms polygon-vertices-length-at-least-4 rotate-polygon-vertices-same-length*
by (*metis (no-types, opaque-lifting) Suc-1 Suc-eq-numeral Suc-le-lessD diff-is-0-eq' eval-nat-numeral(2) gr-implies-not0 length-append-singleton length-butlast length-rotate not-less-eq-eq rotate-polygon-vertices-def*)

let *?n = length vts - 1*
obtain *k* **where** *k: k * ?n > i*
using *len-vts*
by (*metis Suc-1 Suc-le-eq add-0 div-less-iff-less-mult le-add2 less-diff-conv*)
let *?j = k * ?n - i*
have *j-i-n: ?j + i = k * ?n* **using** *k* **by** *simp*

have *rotate-polygon-vertices ?vts' ?j = rotate-polygon-vertices vts (?j + i)*
using *rotation-sum[of vts i ?n]* **by** (*simp add: add commute rotation-sum*)
also have *... = rotate-polygon-vertices vts (k * ?n)* **using** *assms j-i-n* **by** *presburger*

also have ... = vts using polygon-rotate-mod-arb len-vts assms by (metis mult.commute)
 finally show ?thesis using rotation-is-polygon assms by metis
 qed

lemma *rotated-polygon-vertices*:

assumes vts' = rotate-polygon-vertices vts j
 assumes hd vts = last vts
 assumes length vts ≥ 2
 assumes $j \leq i \wedge i < \text{length } vts$
 shows vts ! i = vts' ! (i - j)
 using assms
proof (induct j arbitrary: vts vts')
 case 0
 then show ?case
 by (metis Suc-1 Suc-le-eq diff-is-0-eq diff-zero hd-conv-nth id-apply length-butlast
 linorder-not-le list.size(3) nth-butlast rotate0 rotate-polygon-vertices-def snoc-eq-iff-butlast)
 next
 case (Suc j)
 then have vts' = rotate-polygon-vertices (rotate-polygon-vertices vts 1) j
 by (metis plus-1-eq-Suc rotation-sum)
 moreover have ...!(i - Suc j) = (rotate-polygon-vertices vts 1)!(i - 1)
 using Suc.hyps Suc.premis(3) Suc.premis(4) Suc-1 Suc-diff-le Suc-leD diff-Suc-Suc
 hd-conv-nth length-append-singleton length-butlast length-rotate nth-butlast rotate-polygon-vertices-def
 snoc-eq-iff-butlast zero-less-Suc
 by (smt (z3) One-nat-def Suc.premis(1) Suc.premis(2) Suc-eq-plus1 Suc-le-eq
 arb-rotation-as-single-rotation calculation diff-diff-cancel diff-is-0-eq diff-less-mono
 diff-zero not-less-eq-eq plus-1-eq-Suc rotated-polygon-vertices-helper2)
 moreover have ... = vts!i using rotated-polygon-vertices-helper2
 by (metis Suc.premis(2) Suc.premis(3) Suc.premis(4) add-leD1 le-add-diff-inverse2
 less-diff-conv plus-1-eq-Suc)
 ultimately show ?case
 by presburger
 qed

lemma *polygon-path-image*:

assumes poly-p: polygon p
 assumes p-is-path: p = make-polygonal-path vts
 shows path-image p = p ' {0 ..< 1}
proof -
 have vts-nonempty: vts $\neq []$
 using polygon-at-least-3-vertices[OF poly-p p-is-path]
 by auto
 have at-0: p ' {0} = {pathstart p}
 using p-is-path
 by (metis image-empty image-insert pathstart-def)
 have at-1: p ' {1} = {pathfinish p}
 using p-is-path
 by (simp add: pathfinish-def)
 have same-point: p 0 = p 1

```

using assms unfolding polygon-def closed-path-def using polygon-pathstart[OF
 vts-nonempty p-is-path]
using polygon-pathfinish[OF vts-nonempty p-is-path]
at-0 at-1 by auto
have  $\bigwedge x. x \in p \text{ ' } \{0..1\} \implies x \in p \text{ ' } \{0..<1\}$ 
proof –
  fix x
  assume  $x \in p \text{ ' } \{0..1\}$ 
  then have  $\exists k \in \{0..1\}. p \ k = x$ 
    by auto
  then obtain k where k-prop:  $k \in \{0..1\} \wedge p \ k = x$ 
    by auto
  {assume  $*$ ;  $k < 1$ 
    then have  $\exists k \in \{0..<1\}. p \ k = x$ 
      using k-prop by auto
    } moreover {assume  $*$ ;  $k = 1$ 
    then have  $p \ 0 = x$ 
      using same-point k-prop by auto
    then have  $\exists k \in \{0..<1\}. p \ k = x$ 
      by auto
    }
  ultimately have  $\exists k \in \{0..<1\}. p \ k = x$ 
    using k-prop
    by (metis atLeastAtMost-iff order-less-le)
  then show  $x \in p \text{ ' } \{0..<1\}$ 
    by auto
qed
then show ?thesis
  unfolding path-image-def by auto
qed

lemma polygon-vts-one-rotation:
  fixes p :: R-to-R2
  assumes poly-p: polygon p and
    p-is-path:  $p = \text{make-polygonal-path } vts$  and
    p'-is:  $p' = \text{make-polygonal-path } (\text{rotate-polygon-vertices } vts \ 1)$ 
  shows path-image p = path-image p'
proof –
  let ?rotated-vts = (rotate-polygon-vertices vts 1)
  have  $\text{card } (\text{set } vts) \geq 3$ 
    using polygon-at-least-3-vertices[OF poly-p p-is-path]
    by auto
  then have len-gt-eq3:  $\text{length } vts \geq 3$ 
    using card-length order-trans by blast
  have same-len:  $\text{length } ?rotated-vts = \text{length } vts$ 
    unfolding rotate-polygon-vertices-def using length-rotate
    by (metis One-nat-def Suc-pred card.empty length-append-singleton length-butlast
length-greater-0-conv list.set(1) not-numeral-le-zero p-is-path poly-p polygon-at-least-3-vertices)
  then have len-rotated-gt-eq2:  $\text{length } ?rotated-vts \geq 2$ 

```

```

    using len-gt-eq3 by auto
    have h1:  $\bigwedge x. x \in (\text{path-image } p) \implies x \in \text{path-image } p'$ 
    proof -
      fix x
      assume x  $\in$  (path-image p)
      then have  $\exists k < \text{length } vts - 1. x \in \text{path-image } (\text{linepath } (vts ! k) (vts ! (k + 1)))$ 
      using p-is-path len-gt-eq3 make-polygonal-path-image-property[of vts x]
      by auto
      then obtain k where k-prop:  $k < \text{length } vts - 1 \wedge x \in \text{path-image } (\text{linepath } (vts ! k) (vts ! (k + 1)))$ 
      by auto
      {assume *:  $k = 0$ 
        have vts1:  $vts ! 0 = ?rotated-vts ! (\text{length } ?rotated-vts - 2)$ 
          unfolding rotate-polygon-vertices-def
          using nth-rotate[of length ?rotated-vts - 2 butlast vts 1]
          by (metis (no-types, lifting) * One-nat-def Suc-pred butlast-snoc diff-diff-left
            k-prop length-butlast lessI mod-self nat-1-add-1 nth-butlast plus-1-eq-Suc rotate-polygon-vertices-def
            same-len)
          have (rotate 1 (butlast vts)) ! 0 = vts ! 1
            using nth-rotate[of 0 butlast vts 1] len-gt-eq3
            by (simp add: less-diff-conv mod-if nth-butlast)
          then have vts2:  $vts ! 1 = ?rotated-vts ! (\text{length } ?rotated-vts - 1)$ 
            unfolding rotate-polygon-vertices-def
            by (metis butlast-snoc length-butlast nth-append-length)
          then have path-image (linepath (vts ! k) (vts ! (k + 1)))  $\subseteq$  path-image p'
            using linepaths-subset-make-polygonal-path-image[of vts 0]
            len-rotated-gt-eq2 *
            by (metis (no-types, lifting) One-nat-def Suc-eq-plus1 Suc-pred diff-diff-left
              diff-less k-prop less-numeral-extra(1) linepaths-subset-make-polygonal-path-image nat-1-add-1
              p'-is same-len vts1)
          then have x  $\in$  path-image p'
            using k-prop vts1 vts2
            by auto
        }
      moreover {assume *:  $k > 0$ 
        then have k-minus-prop:  $k - 1 < \text{length } (\text{rotate-polygon-vertices } vts 1) - 1$ 
          using same-len k-prop less-imp-diff-less
          by presburger
        then have vts1:  $vts ! k = ?rotated-vts ! (k - 1)$ 
          using nth-rotate[of k - 1 butlast vts 1] len-gt-eq3
          same-len
          by (metis * One-nat-def Suc-pred butlast-snoc k-prop length-butlast mod-less
            nth-butlast plus-1-eq-Suc rotate-polygon-vertices-def)
        have vts2:  $vts ! (k + 1) = ?rotated-vts ! k$ 
          using nth-rotate[of k butlast vts 1] len-gt-eq3 k-minus-prop
          by (metis (no-types, lifting) * Suc-eq-plus1 Suc-leI butlast-snoc have-wraparound-vertex
            k-prop le-imp-less-Suc length-butlast mod-less mod-self nat-less-le nth-append-length
            nth-butlast p-is-path plus-1-eq-Suc poly-p rotate-polygon-vertices-def same-len)
      }
    }
  
```

```

have path-image (linepath (?rotated-vts ! (k-1)) (?rotated-vts ! k))  $\subseteq$  path-image
p'
  using linepaths-subset-make-polygonal-path-image[OF len-rotated-gt-eq2
k-minus-prop] p'-is
  by (simp add: *)
  then have x  $\in$  path-image p'
  using k-prop vts1 vts2
  by auto
}
ultimately show x  $\in$  path-image p'
by auto
qed
have h2:  $\bigwedge x. x \in (\text{path-image } p') \implies x \in \text{path-image } p$ 
proof -
  fix x
  assume x  $\in$  (path-image p')
  then have  $\exists k < \text{length } ?rotated-vts - 1. x \in \text{path-image } (\text{linepath } (?rotated-vts
! k) (?rotated-vts ! (k + 1)))$ 
  using p'-is len-rotated-gt-eq2 make-polygonal-path-image-property[of ?rotated-vts
x]
  by auto
  then obtain k where k-prop: k < length ?rotated-vts - 1  $\wedge$  x  $\in$  path-image
(linepath (?rotated-vts ! k) (?rotated-vts ! (k + 1)))
  by auto
  {assume *: k = length ?rotated-vts - 2
  have vts1: vts ! 0 = ?rotated-vts ! (length ?rotated-vts - 2)
  unfolding rotate-polygon-vertices-def
  using nth-rotate[of length ?rotated-vts - 2 butlast vts 1]
  by (metis * Suc-diff-Suc Suc-le-eq butlast-snoc k-prop len-rotated-gt-eq2
length-butlast mod-self nat-1-add-1 nth-butlast plus-1-eq-Suc rotate-polygon-vertices-def
same-len zero-less-Suc)
  have (rotate 1 (butlast vts)) ! 0 = vts ! 1
  unfolding rotate-polygon-vertices-def
  using nth-rotate[of 0 butlast vts 1] len-gt-eq3 len-rotated-gt-eq2
  by (metis (no-types, lifting) One-nat-def Suc-le-eq diff-diff-left length-butlast
less-nat-zero-code mod-less not-gr-zero nth-butlast numeral-3-eq-3 plus-1-eq-Suc zero-less-diff)
  then have vts2: ?rotated-vts ! (k+1) = vts ! 1
  unfolding rotate-polygon-vertices-def
  by (metis * Suc-diff-Suc Suc-eq-plus1 Suc-le-eq len-rotated-gt-eq2 length-butlast
length-rotate nat-1-add-1 nth-append-length same-len)
  have path-image (linepath (vts ! 0) (vts ! 1))  $\subseteq$  path-image p
  using linepaths-subset-make-polygonal-path-image[of vts 0]
  len-gt-eq3 * less-diff-conv p-is-path same-len
  by auto
  then have x  $\in$  path-image p
  using * vts1 vts2 k-prop
  by auto
} moreover {assume *: k < length ?rotated-vts - 2
then have vts1: ?rotated-vts ! k = vts ! (k+1)

```

```

using nth-rotate[of k butlast vts 1] len-gt-eq3 *
  same-len
by (smt (z3) Suc-eq-plus1 butlast-snoc diff-diff-left k-prop length-butlast
less-diff-conv mod-less nat-1-add-1 nth-butlast plus-1-eq-Suc rotate-polygon-vertices-def)
have vts2: ?rotated-vts ! (k+1) = vts ! (k+2)
using nth-rotate[of k+1 butlast vts 1] len-gt-eq3 *
by (smt (verit, ccfv-threshold) One-nat-def Suc-le-eq add-Suc-right but-
last-snoc diff-diff-left have-wraparound-vertex len-rotated-gt-eq2 length-butlast less-diff-conv
mod-less mod-self nat-1-add-1 nat-less-le nth-append-length nth-butlast p-is-path
plus-1-eq-Suc poly-p rotate-polygon-vertices-def same-len)
have path-image (linepath (vts ! (k+1)) (vts ! (k + 2)))  $\subseteq$  path-image p
using linepaths-subset-make-polygonal-path-image[of vts k+1]
len-gt-eq3 * less-diff-conv p-is-path same-len
by auto
then have  $x \in$  path-image p
using vts1 vts2 k-prop
by auto
}
ultimately show  $x \in$  path-image p
using k-prop Suc-eq-plus1 add-le-imp-le-diff diff-diff-left len-rotated-gt-eq2
less-diff-conv2 linorder-neqE-nat not-less-eq one-add-one
by linarith
qed
then show ?thesis
using h1 h2 by auto
qed

```

```

lemma polygon-vts-arb-rotation:
fixes p :: R-to-R2
assumes polygon p and
   $p =$  make-polygonal-path vts
shows path-image p = path-image (make-polygonal-path (rotate-polygon-vertices
vts i))
using assms
proof (induct i)
case 0
then show ?case unfolding rotate-polygon-vertices-def
by (metis One-nat-def arb-rotation-as-single-rotation polygon-vts-one-rotation
rotate-polygon-vertices-def rotation-is-polygon)
next
case (Suc i)
let ?p' = make-polygonal-path (rotate-polygon-vertices vts (Suc i))
{assume  $*$ :  $i = 0$ 
have path-image p = path-image ?p'
using Suc polygon-vts-one-rotation[of p vts]
by (simp add: *)
}
moreover {assume  $*$ :  $i > 0$ 
have path-image p = path-image ?p'

```

```

    using polygon-vts-one-rotation arb-rotation-as-single-rotation rotation-is-polygon
      by (metis Suc.hyps Suc.prem1(1) assms(2))
  }
  ultimately show ?case by auto
qed

```

10 Translating a Polygon

lemma *linepath-translation*:

linepath $((\lambda x. x + u) a) ((\lambda x. x + u) b) = (\lambda x. x + u) \circ (\text{linepath } a \ b)$

proof –

let $?l = \text{linepath } ((\lambda x. x + u) a) ((\lambda x. x + u) b)$

let $?l' = (\lambda x. x + u) \circ (\text{linepath } a \ b)$

have $?l \ x = ?l' \ x$ for x

proof –

have $?l \ x = (1 - x) *_R (a + u) + x *_R (b + u)$ **unfolding** *linepath-def* **by** *simp*

also have $\dots = ((1 - x) *_R a + x *_R b) + u$ **by** (*simp add: scaleR-right-distrib*)

also have $\dots = ?l' \ x$ **unfolding** *linepath-def* **by** *simp*

finally show ?thesis .

qed

thus ?thesis **by** *fast*

qed

lemma *make-polygonal-path-translate*:

assumes $\text{length } vts \geq 2$

shows *make-polygonal-path* $(\text{map } (\lambda x. x + u) \ vts) = (\lambda x. x + u) \circ (\text{make-polygonal-path } vts)$

using *assms*

proof(*induct length vts arbitrary: u vts*)

case 0

then show ?case **by** *presburger*

next

case (*Suc n*)

let $?vts' = \text{map } (\lambda x. x + u) \ vts$

let $?p' = \text{make-polygonal-path } ?vts'$

{ assume $\text{Suc } n = 2$

then obtain $a \ b$ where $ab: vts = [a, b]$

by (*metis (no-types, lifting) One-nat-def Suc.hyps(2) Suc-1 Suc-length-conv length-0-conv*)

then have $?vts' = [(\lambda x. x + u) \ a, (\lambda x. x + u) \ b]$ **by** *simp*

then have $?p' = \text{linepath } ((\lambda x. x + u) \ a) ((\lambda x. x + u) \ b)$

using *make-polygonal-path.simps(3)* **by** *presburger*

also have $\dots = (\lambda x. x + u) \circ (\text{linepath } a \ b)$ **using** *linepath-translation* **by** *auto*

also have $\dots = (\lambda x. x + u) \circ (\text{make-polygonal-path } vts)$ **using** *ab* **by** *auto*

finally have ?case .

} moreover

{ assume *: $\text{Suc } n > 2$

```

then obtain  $a\ b\ c\ rest$  where  $abc: vts = a \# b \# c \# rest$ 
  by (metis One-nat-def Suc.hyps(2) Suc-1 Suc-leI Suc-le-length-iff)

let  $?vts\text{-}tl = tl\ vts$ 
let  $?p\text{-}tl = make\text{-}polygonal\text{-}path\ ?vts\text{-}tl$ 
let  $?vts'\text{-}tl = map\ (\lambda x. x + u)\ ?vts\text{-}tl$ 
let  $?p'\text{-}tl = make\text{-}polygonal\text{-}path\ ?vts'\text{-}tl$ 

have  $?vts'\text{-}tl = tl\ ?vts'$  by (simp add: map-tl)
then have  $?p' = (linepath\ (?vts'\!0)\ (?vts'\!1))\ +++\ ?p'\text{-}tl$ 
  using make-polygonal-path.simps(4) abc by force
moreover have  $?p'\text{-}tl = (\lambda x. x + u) \circ (?p\text{-}tl)$  using Suc.hyps(1) Suc.hyps(2)
* by force
  moreover have  $(linepath\ (?vts'\!0)\ (?vts'\!1)) = (\lambda x. x + u) \circ (linepath\ a\ b)$ 
    using abc linepath-translation by auto
  ultimately have  $?case$  by (simp add: abc path-compose-join)
}
ultimately show  $?case$  using Suc by linarith
qed

lemma translation-is-polygon:
  assumes polygon-of p vts
  shows polygon-of (( $\lambda x. x + u$ )  $\circ$  p) (map ( $\lambda x. x + u$ ) vts) (is polygon-of ?p' ?vts')
proof –
  have  $length\ vts \geq 3$ 
  by (metis One-nat-def Suc-eq-plus1 Suc-le-eq add-Suc-right assms nat-less-le numeral-3-eq-3 numeral-Bit0 one-add-one polygon-of-def polygon-vertices-length-at-least-4)
  then have  $*$ :  $?p' = make\text{-}polygonal\text{-}path\ ?vts'$ 
    using make-polygonal-path-translate assms unfolding polygon-of-def by force
  moreover have polygon ?p'
proof –
  have polygonal-path ?p' unfolding polygonal-path-def using  $*$  by simp
  moreover have simple-path ?p'
    using assms unfolding polygon-of-def polygon-def
    using simple-path-translation-eq[of u p]
    by (metis add commute fun.map-cong)
  moreover have closed-path ?p'
proof –
  have  $?p'\ 0 = p\ 0 + u$  by simp
  moreover have  $?p'\ 1 = p\ 1 + u$  by simp
  moreover have  $p\ 0 = p\ 1$ 
    using assms
  unfolding polygon-of-def polygon-def closed-path-def pathstart-def pathfinish-def
  by blast
  moreover have path ?p' using make-polygonal-path-gives-path * by simp
  ultimately show  $?thesis$ 
    unfolding closed-path-def pathstart-def pathfinish-def

```

by *argo*
 qed
 ultimately show *?thesis unfolding polygon-def by blast*
 qed
 ultimately show *?thesis unfolding polygon-of-def by blast*
 qed

11 Misc. properties

lemma *tail-of-loop-free-polygonal-path-is-loop-free:*

assumes *loop-free (make-polygonal-path (x#tail)) (is loop-free ?p) and*
 length tail ≥ 2

shows *loop-free (make-polygonal-path tail) (is loop-free ?p')*

proof –

obtain *y z tail' where tail': tail = y # z # tail'*

by (*metis One-nat-def Suc-1 assms(2) length-Cons list.exhaust-sel list.size(3)*
not-less-eq-eq zero-le)

have *path ?p ∧ path ?p' using make-polygonal-path-gives-path by auto*

have *loop-free ?p using assms unfolding simple-path-def by auto*

moreover have *?p = (linepath x y) +++ ?p'*

using tail' make-polygonal-path.simps(4) by (simp add: tail')

moreover from *calculation have loop-free ?p'*

by (metis make-polygonal-path-gives-path not-loop-free-second-component path-join-path-ends)

ultimately show *?thesis*

using make-polygonal-path-gives-path simple-path-def by blast

qed

lemma *tail-of-simple-polygonal-path-is-simple:*

assumes *simple-path (make-polygonal-path (x#tail)) (is simple-path ?p) and*
 length tail ≥ 2

shows *simple-path (make-polygonal-path tail) (is simple-path ?p')*

using *tail-of-loop-free-polygonal-path-is-loop-free unfolding simple-path-def*

using *assms(1) assms(2) make-polygonal-path-gives-path simple-path-def by blast*

lemma *interior-vtx-in-path-image-interior:*

fixes *vts :: (real²) list*

assumes *x ∈ set (butlast (drop 1 vts))*

shows $\exists t. t \in \{0 < .. < 1\} \wedge (\text{make-polygonal-path } vts) t = x$

using *assms*

proof(*induct vts rule: make-polygonal-path.induct*)

case *1*

then show *?case by simp*

next

case (*2 a*)

then show *?case by simp*

next

case (*3 a b*)

then show *?case by simp*

next

```

case ih: (4 a b c tail')
let ?vts = a # b # c # tail'
let ?tl = b # c # tail'
let ?p = make-polygonal-path ?vts
let ?p-tl = make-polygonal-path ?tl
{ assume  $x \in \text{set } (\text{butlast } (\text{drop } 1 \text{ ?tl}))$ 
  then obtain t' where t':  $t' \in \{0 < .. < 1\} \wedge ?p\text{-tl } t' = x$  using ih by blast
  then have  $?p ((t' + 1) / 2) = x$ 
    unfolding make-polygonal-path.simps joinpaths-def
    by (smt (verit, del-insts) field-sum-of-halves greaterThanLessThan-iff mult-2-right
not-numeral-le-zero zero-le-divide-iff)
    moreover have  $(t' + 1) / 2 \in \{0 < .. < 1\}$  using t' by force
    ultimately have ?case
      by blast
  } moreover
{ assume  $x \notin \text{set } (\text{butlast } (\text{drop } 1 \text{ ?tl}))$ 
  then have  $x = b$ 
  by (metis One-nat-def butlast.simps(2) drop0 drop-Suc-Cons ih.premis list.distinct(1)
set-ConsD)
  then have  $?p (1/2) = x$  unfolding make-polygonal-path.simps joinpaths-def
    by (simp add: linepath-1')
  moreover have  $((1/2)::(\text{real})) \in (\{0 < .. < 1\}::(\text{real set}))$  by simp
  ultimately have ?case by blast
}
ultimately show ?case by auto
qed

```

```

lemma loop-free-polygonal-path-vts-distinct:
  assumes loop-free (make-polygonal-path vts)
  shows distinct (butlast vts)
  using assms
proof(induct vts rule: make-polygonal-path.induct)
  case 1
    then show ?case by simp
  next
    case (2 a)
      then show ?case by simp
  next
    case (3 a b)
      then show ?case by simp
  next
    case ih: (4 a b c tail')
    let ?vts = a # b # c # tail'
    let ?tl = b # c # tail'
    let ?p = make-polygonal-path ?vts
    let ?p-tl = make-polygonal-path ?tl

    have distinct (butlast ?tl)
      using ih tail-of-loop-free-polygonal-path-is-loop-free by simp

```

```

moreover have  $a \notin \text{set } (\text{butlast } ?tl)$ 
proof(rule ccontr)
  assume  $a\text{-in: } \neg a \notin \text{set } (\text{butlast } ?tl)$ 
  then have  $a \in \text{set } (\text{butlast } (\text{drop } 1 \text{ } ?vts))$  by simp
  then obtain  $t$  where  $t: t \in \{0 < .. < 1\} \wedge ?p \ t = a$ 
    using vertices-on-path-image interior-vtx-in-path-image-interior by metis
  then show False
    using ih.premis unfolding simple-path-def loop-free-def
    by (metis atLeastAtMost-iff greaterThanLessThan-iff less-eq-real-def less-numeral-extra(3)
less-numeral-extra(4) list.distinct(1) nth-Cons-0 path-defs(2) polygon-pathstart zero-less-one-class.zero-le-one)
  qed
  ultimately show ?case by simp
qed

```

```

lemma loop-free-polygonal-path-vts-drop1-distinct:
  assumes loop-free (make-polygonal-path vts)
  shows distinct (drop 1 vts)
proof -
  let ?p = make-polygonal-path vts
  let ?last-vts = vts ! ((length vts) - 1)
  have distinct (butlast vts)
  using assms loop-free-polygonal-path-vts-distinct
  by auto
  then have distinct-butlast: distinct (butlast (drop 1 vts))
    by (metis distinct-drop drop-butlast)
  {assume *: length vts > 1
  have len-drop1: length (drop 1 vts) = (length vts) - 1
    using * by simp
  have simp-len:  $1 + ((\text{length } vts) - 2) = (\text{length } vts) - 1$ 
    using * by simp
  then have vts-access:  $vts ! (1 + (\text{length } vts - 2)) = vts ! ((\text{length } vts) - 1)$ 
    by argo
  have drop 1 vts ! ((length vts) - 2) = vts ! (1 + (length vts - 2))
    using * using nth-drop[of 1 vts (length vts) - 2] by auto
  then have ?last-vts = (drop 1 vts) ! ((length vts) - 2)
    using * simp-len vts-access by argo
  then have ?last-vts = (drop 1 vts) ! (length (drop 1 vts) - 1)
    using * len-drop1
    using diff-diff-left nat-1-add-1 by presburger
  then have drop1-is: drop 1 vts = (butlast (drop 1 vts))@[?last-vts]
    using *
  by (metis append-butlast-last-id drop-eq-Nil leD length-butlast nth-append-length)
  have last-vts-not-in: ?last-vts  $\notin \text{set } (\text{butlast } (\text{drop } 1 \text{ } vts))$ 
proof(rule ccontr)
  assume  $a\text{-in: } \neg ?last\text{-vts} \notin \text{set } (\text{butlast } (\text{drop } 1 \text{ } vts))$ 
  then have ?last-vts  $\in \text{set } (\text{butlast } (\text{drop } 1 \text{ } vts))$  by simp
  then obtain  $t$  where  $t: t \in \{0 < .. < 1\} \wedge ?p \ t = ?last\text{-vts}$ 
    using vertices-on-path-image interior-vtx-in-path-image-interior by metis

```

```

have vts ! (length vts - 1) = ?p 1
  using polygon-pathfinish[of vts ?p] *
  by (metis list.size(3) not-one-less-zero pathfinish-def)
then show False
  using t assms unfolding loop-free-def
  by (metis atLeastAtMost-iff greaterThanLessThan-iff leD less-eq-real-def zero-less-one-class.zero-le-one)
qed
have  $\bigwedge b::(\text{real}^2)$  list. distinct b  $\wedge$  a  $\notin$  set b  $\implies$  distinct (b @[a]) for a:: $\text{real}^2$ 
  by simp
then have ?thesis using last-vts-not-in drop1-is distinct-butlast by metis
}
then show ?thesis by force
qed

```

```

lemma simple-polygonal-path-vts-distinct:
  assumes simple-path (make-polygonal-path vts)
  shows distinct (butlast vts)
  using assms loop-free-polygonal-path-vts-distinct
  unfolding simple-path-def
  by blast

```

```

lemma edge-subset-path-image:
  assumes p = make-polygonal-path vts and
    (i::int)  $\in$  {0.. $\langle$ (length vts) - 1 $\rangle$ } and
    x = vts!i and
    y = vts!(i+1)
  shows path-image (linepath x y)  $\subseteq$  path-image p (is ?xy-img  $\subseteq$  ?p-img)
  using assms
proof(induct vts arbitrary: p i rule: make-polygonal-path.induct)
  case 1
  then show ?case by simp
next
  case (2 a)
  then show ?case by simp
next
  case (3 a b)
  then show ?case by (simp add: nth-Cons')
next
  case ih: (4 a b c tl)
  let ?tl = b # c # tl
  let ?p-tl = make-polygonal-path (?tl)
  { assume i = 0
    then have ?case
      by (metis (mono-tags, lifting) ih(2) ih(4) ih(5) Suc-eq-plus1 UnCI list.distinct(1)
        make-polygonal-path.simps(4) nth-Cons-0 nth-Cons-Suc path-image-join pathfinish-linepath polygon-pathstart subsetI)
  } moreover
  { assume i > 0

```

```

then have  $x = ?tl!(i-1)$  by (simp add: ih.premis(3))
moreover have  $y = ?tl!i$  by (simp add: ih.premis(4))
moreover have  $i - 1 \in \{0..<(\text{length } ?tl) - 1\}$  using ih.premis(2) by force
ultimately have  $?xy\text{-img} \subseteq \text{path-image } ?p\text{-tl}$  using ih(1) by (simp add: <0 <
i>)
then have ?case
  unfolding ih(2) make-polygonal-path.simps
  by (smt (verit, ccfv-SIG) UnCI make-polygonal-path.simps(4) make-polygonal-path-gives-path
path-image-join path-join-path-ends subsetI subset-iff)
}
ultimately show ?case by linarith
qed

```

12 Properties of Sublists of Polygonal Path Vertex Lists

lemma *make-polygonal-path-image-append-var:*

```

assumes  $\text{length } vts1 \geq 2$ 
shows  $\text{path-image } (\text{make-polygonal-path } (vts1 @ [v])) = \text{path-image } (\text{make-polygonal-path }
vts1 +++ (\text{linepath } (vts1 ! (\text{length } vts1 - 1)) v))$ 
using assms
proof (induct vts1)
  case Nil
  then show ?case by auto
next
  case (Cons a vts1)
  {assume *:  $\text{length } vts1 = 1$ 
  then obtain b where  $vts1 = [b]$ 
  by (metis Cons-nth-drop-Suc One-nat-def drop0 drop-eq-Nil le-numeral-extra(4)
less-numeral-extra(1))
  then have  $\text{path-image } (\text{make-polygonal-path } ((a \# vts1) @ [v])) =$ 
 $\text{path-image } (\text{make-polygonal-path } (a \# vts1) +++ \text{linepath } ((a \# vts1) !$ 
 $(\text{length } (a \# vts1) - 1)) v)$ 
  using make-polygonal-path.simps
  by simp
} moreover {assume *:  $\text{length } vts1 > 1$ 
  then obtain b c vts1' where  $vts1 = b \# c \# vts1'$ 
  by (metis One-nat-def length-0-conv length-Cons less-numeral-extra(4) not-one-less-zero
remdups-adj.cases)
  then have  $h1: \text{make-polygonal-path } ((a \# vts1) @ [v]) = (\text{linepath } a b) +++$ 
 $(\text{make-polygonal-path } (vts1 @ [v]))$ 
  using make-polygonal-path.simps(4)
  by auto
  have  $\text{path-image } (\text{make-polygonal-path } (vts1 @ [v])) =$ 
 $\text{path-image } (\text{make-polygonal-path } vts1 +++ \text{linepath } (vts1 ! (\text{length } vts1 - 1))$ 
 $v)$ 
  using * Cons by auto
  then have  $\text{path-image } (\text{make-polygonal-path } ((a \# vts1) @ [v])) =$ 

```

```

  path-image (make-polygonal-path (a # vts1) +++ linepath ((a # vts1) ! (length
(a # vts1) - 1)) v)
  using h1
  by (metis (no-types, lifting) Cons.prem1 Suc-1 Suc-le-eq Un-assoc ‹vts1 = b # c
# vts1 ‹ add-diff-cancel-left' append-Cons length-Cons list.discI make-polygonal-path.simps(4)
nth-Cons-0 nth-Cons-pos path-image-join pathfinish-linepath pathstart-linepath plus-1-eq-Suc
polygon-pathfinish polygon-pathstart zero-less-diff)
}
ultimately show ?case
  by (metis Cons.prem1 Suc-1 add-diff-cancel-left' le-neq-implies-less length-Cons
not-less-eq plus-1-eq-Suc)
qed

```

lemma *make-polygonal-path-image-append-helper:*

```

  assumes length vts1 ≥ 1 ∧ length vts2 ≥ 1
  shows path-image (make-polygonal-path (vts1 @ [v] @ [v] @ vts2)) = path-image
(make-polygonal-path (vts1 @ [v] @ vts2))
  using assms
proof (induct vts1)
  case Nil
  then show ?case by auto
next
  case (Cons a vts1)
  { assume *: length vts1 = 0
  have path-image (make-polygonal-path ([a] @ [v] @ vts2)) =
    path-image ((linepath a v) +++ make-polygonal-path (v # vts2))
  using make-polygonal-path.simps
  by (metis Cons.prem1 One-nat-def append-Cons append-Nil append-take-drop-id
linorder-not-le list.distinct(1) list.exhaust not-less-eq-eq take-hd-drop)
  then have path-image (make-polygonal-path ([a] @ [v] @ vts2)) =
    path-image (linepath a v) ∪ path-image (make-polygonal-path (v # vts2))
  by (metis list.discI nth-Cons-0 path-image-join pathfinish-linepath polygon-pathstart)
  have image-helper1: path-image (make-polygonal-path ([a] @ [v] @ [v] @ vts2))
= path-image (linepath a v +++ make-polygonal-path (v # v # vts2))
  by simp
  have path-image (make-polygonal-path (v # v # vts2)) = path-image ((linepath
v v) +++ make-polygonal-path (v # vts2))
  using make-polygonal-path.simps
  by (metis Cons.prem1 One-nat-def append-Cons append-Nil append-take-drop-id
linorder-not-le list.distinct(1) list.exhaust not-less-eq-eq take-hd-drop)
  moreover have ... = path-image (linepath v v) ∪ path-image (make-polygonal-path
(v # vts2))
  by (metis list.discI nth-Cons-0 path-image-join pathfinish-linepath poly-
gon-pathstart)
  ultimately have image-helper2: path-image (make-polygonal-path (v # v #
vts2)) = {v} ∪ path-image (make-polygonal-path (v # vts2))
  by auto
  have v ∈ path-image (make-polygonal-path (v # vts2))
  using vertices-on-path-image by fastforce

```

```

then have path-image (make-polygonal-path ([a] @ [v] @ [v] @ vts2)) =
path-image (make-polygonal-path ([a] @ [v] @ vts2))
using image-helper1 image-helper2
by (metis ‹path-image (make-polygonal-path ([a] @ [v] @ vts2)) = path-image
(linepath a v) ∪ path-image (make-polygonal-path (v # vts2))› insert-absorb in-
sert-is-Un list.simps(3) nth-Cons-0 path-image-join pathfinish-linepath polygon-pathstart)
}
moreover {assume *: length vts1 > 0
then have ind-hyp: path-image (make-polygonal-path (vts1 @ [v] @ [v] @ vts2))
=
path-image (make-polygonal-path (vts1 @ [v] @ vts2))
using Cons.hyps Cons.prem by linarith
obtain b vts3 where vts1-is: vts1 = b # vts3
using *
by (metis * Cons-nth-drop-Suc drop0)
then have path-image1: path-image (make-polygonal-path ((a # vts1) @ [v] @
[v] @ vts2)) =
path-image ((linepath a b) +++ make-polygonal-path (vts1 @ [v] @ [v] @
vts2))
by (smt (verit, best) Cons.prem Nil-is-append-conv append-Cons length-greater-0-conv
less-numeral-extra(1) list.inject make-polygonal-path.elims order-less-le-trans)
obtain c d where bcd: vts1 @ [v] @ vts2 = b # c # d
using vts1-is
by (metis append-Cons append-Nil neq-Nil-conv)
have path-image2: path-image (make-polygonal-path ((a # vts1) @ [v] @ vts2))
= path-image ((linepath a b) +++ make-polygonal-path (vts1 @ [v] @ vts2))
using make-polygonal-path.simps bcd
by auto
have path-image (make-polygonal-path ((a # vts1) @ [v] @ [v] @ vts2)) =
path-image (make-polygonal-path ((a # vts1) @ [v] @ vts2))
using ind-hyp path-image1 path-image2
by (smt (verit, del-insts) Nil-is-append-conv append-Cons nth-Cons-0 path-image-join
pathfinish-linepath polygon-pathstart vts1-is)
}
ultimately show ?case
using Cons.prem
by blast
qed

```

```

lemma make-polygonal-path-image-append:
assumes length vts1 ≥ 2 ∧ length vts2 ≥ 2
shows path-image (make-polygonal-path (vts1 @ vts2)) = path-image (make-polygonal-path
vts1 +++ (linepath (vts1 ! (length vts1 - 1)) (vts2 ! 0)) +++ make-polygonal-path
vts2)
using assms
proof (induct vts1)
case Nil
then show ?case
by simp

```

```

next
  case (Cons a vts1)
  {assume *: length vts1 = 1
   then obtain b where vts1-is: vts1 = [b]
   by (metis Cons-nth-drop-Suc One-nat-def drop0 drop-eq-Nil le-numeral-extra(4)
less-numeral-extra(1))
   then have make-polygonal-path ((a # vts1) @ vts2) = make-polygonal-path (a
# b # vts2)
   by simp
   then have make-polygonal-path ((a # vts1) @ vts2) = (linepath a b) +++
(make-polygonal-path (b # vts2))
   by (metis Cons.premis length-0-conv make-polygonal-path.simps(4) neq-Nil-conv
not-numeral-le-zero)
   then have make-polygonal-path ((a # vts1) @ vts2) = make-polygonal-path
(a # vts1) +++ (make-polygonal-path (b # vts2))
   using vts1-is make-polygonal-path.simps(3)
   by simp
   then have make-polygonal-path ((a # vts1) @ vts2) = make-polygonal-path
(a # vts1) +++ linepath b (vts2 ! 0) +++ make-polygonal-path vts2
   using Cons.premis
   by (smt (verit, ccfv-SIG) * Suc-1 add-diff-cancel-left' diff-is-0-eq' length-greater-0-conv
list.size(4) make-polygonal-path.elims make-polygonal-path.simps(4) nth-Cons-0 or-
der-less-le-trans plus-1-eq-Suc pos2 vts1-is zero-neq-one)
   then have make-polygonal-path ((a # vts1) @ vts2) =
make-polygonal-path (a # vts1) +++
linepath ((a # vts1) ! (length (a # vts1) - 1)) (vts2 ! 0) +++ make-polygonal-path
vts2
   using vts1-is
   by simp
} moreover {assume *: length vts1 > 1
 then obtain b c vts1' where vts1 = b # c # vts1'
 by (metis One-nat-def length-0-conv length-Cons less-numeral-extra(4) not-one-less-zero
remdups-adj.cases)
 then have h1: make-polygonal-path ((a # vts1) @ vts2) = (linepath a b) +++
(make-polygonal-path (vts1 @ vts2))
   using make-polygonal-path.simps(4)
   by auto
   have ind-h: path-image (make-polygonal-path (vts1 @ vts2)) =
path-image (make-polygonal-path vts1 +++
linepath (vts1 ! (length vts1 - 1)) (vts2 ! 0) +++ make-polygonal-path vts2)
   using Cons *
   by linarith
 then have path-image (make-polygonal-path ((a # vts1) @ vts2)) = path-image
((linepath a b) ∪ path-image((make-polygonal-path vts1 +++
linepath (vts1 ! (length vts1 - 1)) (vts2 ! 0) +++ make-polygonal-path vts2))
   using h1
   by (metis (mono-tags, lifting) * Nil-is-append-conv ⟨vts1 = b # c # vts1'⟩ ap-
pend-Cons length-greater-0-conv linordered-nonzero-semiring-class.zero-le-one nth-Cons-0
order-le-less-trans path-image-join pathfinish-linepath polygon-pathstart)

```

then have $\text{path-image (make-polygonal-path ((a \# vts1) @ vts2))} = (\text{path-image (linepath a b)} \cup \text{path-image (make-polygonal-path vts1)}) \cup \text{path-image}((\text{linepath (vts1 ! (length vts1 - 1)) (vts2 ! 0)} \text{ +++ make-polygonal-path vts2}))$
by (*metis (no-types, lifting) * Un-assoc length-greater-0-conv order-le-less-trans path-image-join pathstart-join pathstart-linepath polygon-pathfinish zero-less-one-class.zero-le-one*)
then have $\text{image-helper: path-image (make-polygonal-path ((a \# vts1) @ vts2))} = (\text{path-image (make-polygonal-path (a \# vts1))}) \cup \text{path-image}((\text{linepath (vts1 ! (length vts1 - 1)) (vts2 ! 0)} \text{ +++ make-polygonal-path vts2}))$
by (*metis (no-types, lifting) * \langle vts1 = b \# c \# vts1' \rangle length-greater-0-conv make-polygonal-path.simps(4) nth-Cons-0 order-le-less-trans path-image-join pathfinish-linepath polygon-pathstart zero-less-one-class.zero-le-one*)
have $vts1 ! (length vts1 - 1) = (a \# vts1) ! (length (a \# vts1) - 1)$
using *Cons.prem*s
by (*simp add: Suc-le-eq*)
then have $\text{path-image (make-polygonal-path ((a \# vts1) @ vts2))} = \text{path-image (make-polygonal-path (a \# vts1) \text{ +++ linepath ((a \# vts1) ! (length (a \# vts1) - 1)) (vts2 ! 0) \text{ +++ make-polygonal-path vts2})}$
using *image-helper*
by (*metis (no-types, lifting) Cons.prem*s *length-greater-0-conv order-less-le-trans path-image-join pathstart-join pathstart-linepath polygon-pathfinish pos2*)
}
ultimately show *?case using Cons.prem*s
by *fastforce*
qed

lemma *make-polygonal-path-image-append-alt:*

assumes $p = \text{make-polygonal-path vts}$
assumes $p1 = \text{make-polygonal-path vts1}$
assumes $p2 = \text{make-polygonal-path vts2}$
assumes $\text{last vts1} = \text{hd vts2}$
assumes $\text{length vts1} \geq 2 \wedge \text{length vts2} \geq 2$
assumes $vts = vts1 @ (\text{tl vts2})$
shows $\text{path-image } p = \text{path-image } (p1 \text{ +++ } p2)$

proof –

have $\text{path-image } p = \text{path-image } p1 \cup \text{path-image } p2$
by (*smt (z3) Nitpick.size-list-simp(2) One-nat-def Suc-1 assms diff-Suc-1 last-conv-nth length-greater-0-conv list.collapse list.sel(3) make-polygonal-path.elims make-polygonal-path.simps(3) make-polygonal-path-image-append make-polygonal-path-image-append-var nat-less-le not-less-eq-eq nth-Cons-0 order-less-le-trans path-image-join polygon-pathfinish polygon-pathstart pos2 length-Cons length-tl path-image-cons-union pathfinish-linepath pathstart-join sup.absorb-iff1 sup.absorb-iff2*)
thus *?thesis*
by (*metis assms(2) assms(3) assms(4) assms(5) hd-conv-nth last-conv-nth length-greater-0-conv order-less-le-trans path-image-join polygon-pathfinish polygon-pathstart pos2*)

qed

lemma *cont-incr-interval-image*:

fixes $f :: \text{real} \Rightarrow \text{real}$

assumes $a \leq b$

assumes *continuous-on* $\{a..b\}$ f

assumes $\forall x \in \{a..b\}. \forall y \in \{a..b\}. x \leq y \longrightarrow f x \leq f y$

shows $f'\{a..b\} = \{f a..f b\}$

proof –

have $f'\{a..b\} \subseteq \{f a..f b\}$

proof(*rule subsetI*)

fix x

assume $x \in f'\{a..b\}$

then obtain t **where** $t \in \{a..b\} \wedge f t = x$ **by** *blast*

moreover then have $a \leq t \wedge t \leq b$ **by** *presburger*

ultimately show $x \in \{f a..f b\}$ **using** *assms(3)* **by** *auto*

qed

moreover have $\{f a..f b\} \subseteq f'\{a..b\}$

proof –

obtain $c d$ **where** $f'\{a..b\} = \{c..d\}$ **using** *continuous-image-closed-interval*
assms **by** *meson*

moreover then have $f a \in \{c..d\}$ **using** *assms(1)* **by** *auto*

moreover have $f b \in \{c..d\}$ **using** *assms(1)* *calculation* **by** *auto*

moreover have $\{f a..f b\} \subseteq \{c..d\}$ **using** *calculation* **by** *simp*

ultimately show *?thesis* **by** *presburger*

qed

ultimately show *?thesis* **by** *blast*

qed

lemma *two-x-minus-one-image*:

assumes $f = (\lambda x::\text{real}. 2*x - 1)$

assumes $a \leq b$

shows $f'\{a..b\} = \{f a..f b\}$

proof –

have *continuous-on* $\{a..b\}$ f

proof –

have *continuous-on* $\{a..b\}$ $(\lambda x::\text{real}. x)$ **by** *simp*

then have *continuous-on* $\{a..b\}$ $(\lambda x::\text{real}. 2*x)$ **using** *continuous-on-mult-const*
by *blast*

thus *continuous-on* $\{a..b\}$ f

unfolding *assms* **using** *continuous-on-translation-eq*[of $\{a..b\} - 1$ $(\lambda x::\text{real}.$
 $2*x)$] **by** *auto*

qed

thus *?thesis* **using** *cont-incr-interval-image* *assms* **by** *force*

qed

lemma *vts-split-path-image*:

assumes $p = \text{make-polygonal-path } vts$

assumes $p1 = \text{make-polygonal-path } vts1$

```

assumes  $p2 = \text{make-polygonal-path } vts2$ 
assumes  $vts1 = \text{take } i \text{ } vts$ 
assumes  $vts2 = \text{drop } (i-1) \text{ } vts$ 
assumes  $n = \text{length } vts$ 
assumes  $1 \leq i \wedge i < n$ 
assumes  $x = (2^{(i-1)} - 1) / (2^{(i-1)})$ 
shows  $\text{path-image } p1 = p\{0..x\} \wedge \text{path-image } p2 = p\{x..1\}$ 
using  $assms$ 
proof(induct  $i$  arbitrary:  $p$   $p1$   $p2$   $vts$   $vts1$   $vts2$   $n$   $x$ )
  case  $0$ 
    then show  $?case$  by  $linarith$ 
next
  case ( $Suc\ i$ )
    { assume  $*$ :  $Suc\ i = 1$ 
      then obtain  $a$  where  $a$ :  $vts1 = [a]$ 
        using  $Suc.prem$ s
          by (metis  $One\text{-nat-def}$   $gr\text{-implies-not0}$   $list.collapse$   $list.size(3)$   $take\text{-eq-}Nil$ 
             $take\text{-tl}\ zero\text{-neq-one}$ )
          moreover have  $vts2 = vts$  using  $*$   $Suc.prem$ s by  $force$ 
          ultimately have  $p1 = \text{linepath } a\ a \wedge p2 = p$ 
            using  $Suc.prem$ s  $make\text{-polygonal-path.sims}$  by  $meson$ 
          moreover have  $x = 0$  using  $Suc.prem$ s  $*$  by  $simp$ 
          moreover have  $\text{path-image } p1 = \{a\}$  using  $calculation$  by  $simp$ 
          moreover have  $p\{0..0\} = \{p\ 0\}$  by  $auto$ 
          moreover then have  $p\{0..0\} = \{a\}$  using  $Suc.prem$ s
            by (metis  $a\ gr0\text{-conv-Suc}$   $list.discI$   $nth\text{-Cons-}0$   $nth\text{-take}$   $pathstart\text{-def}$   $poly\text{-gon-pathstart}$ 
               $take\text{-eq-}Nil$ )
            moreover have  $\text{path-image } p1 = p\{0..x\}$  using  $calculation$  by  $presburger$ 
            moreover have  $\text{path-image } p2 = p\{x..1\}$  using  $calculation$  unfolding  $path\text{-image-def}$ 
by  $fast$ 
          ultimately have  $?case$  by  $blast$ 
        } moreover
      { assume  $*$ :  $Suc\ i > 1$ 

        let  $?a = vts!0$ 
        let  $?b = vts!1$ 
        let  $?l = \text{linepath } ?a\ ?b$ 
        let  $?L = \text{path-image } ?l$ 
        let  $?tl = \text{tl } vts$ 
        let  $?vts1' = \text{take } i \text{ } ?tl$ 
        let  $?vts2' = \text{drop } (i-1) \text{ } ?tl$ 
        let  $?p' = \text{make-polygonal-path } ?tl$ 
        let  $?p1' = \text{make-polygonal-path } ?vts1'$ 
        let  $?p2' = \text{make-polygonal-path } ?vts2'$ 
        let  $?x' = ((2::real)^{(i-1)} - 1) / (2^{(i-1)})$ 
        let  $?P1' = \text{path-image } ?p1'$ 
        let  $?P2' = \text{path-image } ?p2'$ 

        have  $i$ :  $1 \leq i \wedge i < \text{length } ?tl$ 

```

```

using Suc.prems * by (metis Suc-eq-plus1 length-tl less-Suc-eq-le less-diff-conv)
then have ih:  $?P1' = ?p'\{0..?x'\} \wedge ?P2' = ?p'\{?x'..1\}$ 
using Suc.hyps[of  $?p' ?tl ?p1' ?vts1' ?p2' ?vts2' length ?tl ?x'$ ] by presburger

let  $?f = \lambda x::real. 2*x - 1$ 

have fx:  $?f x = ?x'$ 
by (metis i Suc.prems(8) bounding-interval-helper1 diff-Suc-1 summation-helper)

moreover have fhalf:  $?f (1/2) = 0$  by simp
moreover have f1:  $?f 1 = 1$  by simp
ultimately have f:  $?f\{x..1\} = \{?x'..1\} \wedge ?f\{1/2..x\} = \{0..?x'\}$ 
using two-x-minus-one-image by auto
have x:  $1/2 \leq x \wedge x \leq 1$ 
by (smt (verit) divide-le-eq-1-pos divide-nonneg-nonneg fhalf fx two-realpow-ge-one)

have  $n \geq 3$  using Suc.prems * by linarith
then have p:  $p = ?l +++ ?p'$ 
proof -
  have f1:  $\forall vs. (vs::(real, 2) \text{ vec list}) \neq [] \vee \neg 1 < \text{Suc } (\text{length } vs)$ 
    by simp
  have  $1 < \text{Suc } n$ 
    using Suc.prems(7) by linarith
  then show thesis
    by (smt (verit) f1 Suc-le-lessD i One-nat-def Suc.prems(6) Suc.prems(7)
Suc-less-eq  $\langle p = \text{make-polygonal-path } vts \rangle$  hd-conv-nth length-Cons length-tl less-Suc-eq
list.collapse list.exhaust make-polygonal-path.simps(4) nth-Cons-Suc zero-order(3))

qed
have p-to-p':  $\forall y \geq 1/2. p y = (?p' \circ ?f) y$ 
proof clarify
  fix y :: real
  assume *:  $y \geq 1/2$ 
  { assume **:  $y = 1/2$ 
    then have  $p y = ?b$ 
      by (smt (verit) fhalf joinpaths-def linepath-1' p)
    moreover have  $?f y = 0$  using ** by simp
    moreover have  $?p' 0 = ?b$ 
      by (metis i One-nat-def Suc.prems(6) length-greater-0-conv length-tl
list.size(3) nth-tl pathstart-def polygon-pathstart zero-order(3))
    ultimately have  $p y = (?p' \circ ?f) y$  by simp
  } moreover
  { assume **:  $y > 1/2$ 
    then have  $p y = ?p' (?f y)$  unfolding p joinpaths-def by simp
    then have  $p y = (?p' \circ ?f) y$  by force
  }
  ultimately show  $p y = (?p' \circ ?f) y$  using * by fastforce
qed

```

```

have {0..x} = {0..1/2} ∪ {1/2..x} using x by (simp add: ivl-disj-un-two-touch(4))
then have p'{0..x} = p'{0..1/2} ∪ p'{1/2..x} by blast
also have ... = ?L ∪ p'{1/2..x}
proof-
  have ?L ⊆ p'{0..1/2}
  proof(rule subsetI)
    fix a
    assume *: a ∈ ?L
    then obtain t where t: t ∈ {0..1} ∧ ?l t = a unfolding path-image-def
  by blast
    then have p (t/2) = a unfolding p joinpaths-def by auto
    moreover have t/2 ∈ {0..1/2} using t by simp
    ultimately show a ∈ p'{0..1/2} by blast
  qed
  moreover have p'{0..1/2} ⊆ ?L
  proof(rule subsetI)
    fix a
    assume *: a ∈ p'{0..1/2}
    then obtain t where t ∈ {0..1/2} ∧ p t = a by blast
    moreover then have ?l (2*t) = p t unfolding p joinpaths-def by presburger
    moreover have 2*t ∈ {0..1} using calculation by simp
    ultimately show a ∈ ?L unfolding path-image-def by auto
  qed
  ultimately have ?L = p'{0..1/2} by blast
  thus ?thesis by presburger
qed
also have ... = ?L ∪ (?p' ∘ ?f){1/2..x} using p-to-p' by simp
also have ... = ?L ∪ ?p'{0..?x'} using f by (metis image-comp)
also have ... = ?L ∪ ?P1' using ih by blast
also have ... = path-image p1
proof-
  have take i (tl vts) ≠ [] by (metis i less-zeroE list.size(3) not-one-le-zero
take-eq-Nil2)
  thus ?thesis using path-image-cons-union[of p1 vts1 ?p1' ?vts1' ?a ?b]
  by (metis * Nitpick.size-list-simp(2) One-nat-def Suc.prem(2) Suc.prem(4)
Suc.prem(6) Suc.prem(7) bot-nat-0.extremum-strict hd-conv-nth length-greater-0-conv
nth-take nth-tl take-Suc take-tl)
qed
finally have 1: path-image p1 = p'{0..x} by argo

have p'{x..1} = (?p' ∘ ?f){x..1} using p-to-p' x by simp
also have ... = ?p'{?x'..1} using f by (metis image-comp)
also have ... = ?P2' using ih by presburger
also have ... = path-image p2
using path-image-cons-union
by (metis Suc.prem(3) Suc.prem(5) diff-Suc-1 drop-Suc gr0-implies-Suc i
linorder-neqE-nat not-less-zero not-one-le-zero)
finally have 2: path-image p2 = p'{x..1} by argo

```

```

    have ?case using 1 2 by fast
  }
  ultimately show ?case using Suc.prem1 by linarith
qed

```

lemma *drop-i-is-loop-free:*

```

  fixes vts :: (real^2) list
  assumes m = length vts
  assumes i ≤ m - 2
  assumes vts' = drop i vts
  assumes p = make-polygonal-path vts
  assumes p' = make-polygonal-path vts'
  assumes loop-free p
  shows loop-free p'
  using assms
proof(induct i arbitrary: vts' p')
  case 0
  then show ?case by simp
next
  case (Suc i)

  let ?vts'' = drop i vts
  let ?p'' = make-polygonal-path ?vts''
  have ih: loop-free ?p''
    using Suc.hyps Suc.prem1(2) Suc.prem1(6) Suc-leD assms(1) assms(4) by
blast

```

```

  obtain a b where ab: ?vts'' = a # vts' ! 1 # (drop 2 vts')
  by (metis Cons-nth-drop-Suc Suc.prem1(3) constant-linepath-is-not-loop-free
drop-eq-Nil ih linorder-not-less make-polygonal-path.simps(1))
  then have ?vts'' = a # b # (vts' ! 1) # (drop 2 vts')
  by (smt (verit, ccfv-threshold) Cons-nth-drop-Suc Suc.prem1(2) Suc.prem1(3)
Suc-1 Suc-diff-Suc Suc-le-eq assms(1) diff-Suc-1 diff-is-0-eq drop-drop le-add-diff-inverse
length-drop nat-le-linear not-less-eq-eq zero-less-Suc)
  then have ?p'' = (linepath a b) +++ p'
  using make-polygonal-path.simps(4)[of a b vts' ! 1 drop 2 vts'] Suc.prem1 by
(simp add: ab)
  moreover have pathfinish (linepath a b) = pathstart p'
  using Suc.prem1 ab
  by (metis constant-linepath-is-not-loop-free ih make-polygonal-path.simps(2)
pathfinish-linepath polygon-pathstart)
  ultimately have arc p' using simple-path-joinE
  by (metis ih make-polygonal-path-gives-path simple-path-def)
  then show ?case using arc-imp-simple-path simple-path-def by blast
qed

```

lemma *joinpaths-tl-transform:*

```

  assumes f = (λx::real. 2*x - 1)

```

assumes $\text{pathfinish } g1 = \text{pathstart } g2$
assumes $p = g1 \text{ +++ } g2$
assumes $x \geq 1/2$
shows $p \ x = g2 \ (f \ x)$
proof –
{ **assume** $x = 1/2$
moreover then have $f \ x = 0$ **using** *assms* **by** *fastforce*
ultimately have $p \ x = \text{pathfinish } g1 \wedge g2 \ (f \ x) = \text{pathfinish } g1$
using *assms* **unfolding** *pathfinish-def pathstart-def joinpaths-def* **by** *force*
then have $p \ x = g2 \ (f \ x)$ **using** *assms* **unfolding** *joinpaths-def* **by** *simp*
} **moreover**
{ **assume** $x > 1/2$
then have $p \ x = g2 \ (f \ x)$ **using** *assms* **unfolding** *joinpaths-def* **by** *simp*
}
ultimately show $p \ x = g2 \ (f \ x)$ **using** *assms* **by** *fastforce*
qed

lemma *joinpaths-tl-image-transform*:
assumes $f = (\lambda x::\text{real}. 2*x - 1)$
assumes $\text{pathfinish } g1 = \text{pathstart } g2$
assumes $p = g1 \text{ +++ } g2$
assumes $1/2 \leq a \wedge a \leq b$
shows $p \ \{a..b\} = g2 \ \{f \ a..f \ b\}$
proof –
have $\forall x \in \{a..b\}. p \ x = g2 \ (f \ x)$ **using** *assms* *joinpaths-tl-transform*[of *f* *g1* *g2*
p] **by** *force*
then have $p \ \{a..b\} = (g2 \circ f) \ \{a..b\}$ **by** *simp*
also have $\dots = g2 \ \{f \ a..f \ b\}$ **using** *two-x-minus-one-image* **by** (*metis* *assms*(1,4)
image-comp)
finally show *?thesis* .
qed

lemma *vts-sublist-path-image*:
assumes $p = \text{make-polygonal-path } vts$
assumes $p' = \text{make-polygonal-path } vts'$
assumes $vts' = \text{take } j \ (\text{drop } i \ vts)$
assumes $m = \text{length } vts$
assumes $n = \text{length } vts'$
assumes $k = i + j$
assumes $k \leq m - 1 \wedge 2 \leq j$
assumes $x1 = (2^i - 1)/(2^i)$
assumes $x2 = (2^{(k-1)} - 1)/(2^{(k-1)})$
shows $\text{path-image } p' = p \ \{x1..x2\}$
using *assms*
proof (*induct* *i* *arbitrary*: *vts* *p* *p'* *vts'* *m* *k* *x1* *x2*)
case 0
then show *?case* **using** *vts-split-path-image*[of *p* *drop* 0 *vts* *p'* *vts'* - - *j* *m* *x2*]
by (*metis* (*no-types*, *opaque-lifting*) *Suc-diff-le* *add-0* *cancel-comm-monoid-add-class*.*diff-cancel*
diff-is-0-eq *div-by-1* *drop.simps*(1) *drop-0* *le-add-diff-inverse* *length-drop* *less-one*)

```

linorder-not-le plus-1-eq-Suc pos2 power.simps(1))
next
  case (Suc i)

  let ?vts-tl = tl vts
  let ?vts-tl' = take j (drop i ?vts-tl)
  let ?p-tl = make-polygonal-path ?vts-tl
  let ?m' = m-1
  let ?k' = i+j
  let ?x1' = (2i - 1)/(2i)
  let ?x2' = (2?k'-1 - 1)/(2?k'-1)
  let ?f = λx. 2*x - 1

  have vts' = ?vts-tl' using Suc.prem by (metis drop-Suc)
  then have p' = make-polygonal-path ?vts-tl' using Suc.prem by argo
  then have ih: path-image p' = ?p-tl' { ?x1' .. ?x2' }
    using Suc.hyps [of ?p-tl ?vts-tl p' ?vts-tl' ?m' ?k' ?x1' ?x2' ] Suc.prem
    by (smt (verit, ccfv-SIG) Suc-eq-plus1 add-diff-cancel-right' add-leD1 diff-diff-left
diff-is-0-eq drop-Suc le-add-diff-inverse length-tl linorder-not-le not-add-less2)

  let ?a = vts!0
  let ?b = vts!1
  let ?l = linepath ?a ?b
  have p: p = ?l +++ ?p-tl
  proof-
    have length vts ≥ 3 using Suc.prem by linarith
    then obtain c w where vts = ?a # ?b # c # w
      by (metis Cons-nth-drop-Suc One-nat-def Suc-le-eq drop0 numeral-3-eq-3
order-less-le)
    thus ?thesis
      using Suc.prem make-polygonal-path.simps(4) [of ?a ?b c w] by (metis
list.sel(3))
  qed
  moreover have x1 ≥ 1/2 using Suc.prem by (simp add: plus-1-eq-Suc)
  moreover have x2 ≥ x1
    using Suc.prem
    by (smt (verit, best) Nat.diff-add-assoc2 One-nat-def add-Suc-shift add-diff-cancel-left'
add-mono-thms-linordered-semiring(2) diff-add-cancel dual-order.trans group-cancel.rule0
numeral-One one-le-numeral one-le-power plus-1-eq-Suc power-increasing real-shrink-le
trans-le-add2)
  moreover have pathfinish ?l = pathstart ?p-tl
    by (metis One-nat-def Suc.prem(4) Suc.prem(6) Suc.prem(7) Suc-neq-Zero
add-is-0 diff-is-0-eq' diff-zero length-tl linorder-not-less list.size(3) nth-tl pathfin-
ish-linepath polygon-pathstart)
  ultimately have p {x1..x2} = ?p-tl {?f x1..?f x2}
    using joinpaths-tl-image-transform [of ?f ?l ?p-tl p x1 x2] by presburger
  also have ... = ?p-tl {?x1'..?x2'}
    by (metis (no-types, lifting) Nat.add-diff-assoc Suc.prem(6-9) add commute
add-leD1 bounding-interval-helper1 diff-Suc-1 le-add2 nat-1-add-1 plus-1-eq-Suc sum-

```

```

mation-helper)
  also have ... = path-image p' using ih by blast
  finally show ?case by argo
qed

lemma one-append-simple-path:
  fixes vts :: (real^2) list
  assumes vts = vts' @ [z]
  assumes n = length vts
  assumes n ≥ 3
  assumes p = make-polygonal-path vts
  assumes p' = make-polygonal-path vts'
  assumes simple-path p
  shows simple-path p'
  using assms
proof(induct n arbitrary: vts vts' p p')
  case 0
  then show ?case by linarith
next
  case (Suc n)
  { assume *: Suc n = 3
    then obtain a b c where abc: vts = [a, b, c] ∧ vts' = [a, b]
      using Suc.premis
      by (smt (z3) Suc-le-length-iff Suc-length-conv append-Cons diff-Suc-1 drop0
length-0-conv length-append-singleton numeral-3-eq-3)
    then have p' = linepath a b
      by (simp add: Suc.premis(5))
    moreover have a ≠ b using loop-free-polygonal-path-vts-distinct Suc.premis
      by (metis abc butlast-snoc distinct-length-2-or-more simple-path-def)
    ultimately have ?case by blast
  } moreover
  { assume *: Suc n > 3
    then obtain a b tl' where ab: vts' = a # tl' ∧ b = tl'!0 using Suc.premis
      by (metis Suc-le-length-iff Suc-le-mono length-append-singleton numeral-3-eq-3)
    moreover then have p = make-polygonal-path (a # (tl' @ [z])) using Suc.premis
  }
by auto
  moreover then have p: p = linepath a b +++ make-polygonal-path (tl' @ [z])
    using make-polygonal-path.simps ab
    by (smt (verit, ccfv-threshold) * Cons-nth-drop-Suc One-nat-def Suc.premis(1)
Suc.premis(2) Suc-1 Suc-less-eq append-Cons drop0 length-Cons length-append-singleton
length-greater-0-conv list.size(3) not-numeral-less-one numeral-3-eq-3)
  moreover then have simple-path ... using Suc.premis by meson
  ultimately have pre-ih: simple-path (make-polygonal-path (tl' @ [z]))
    using Suc.premis(1) Suc.premis(2) Suc.premis(3) ab tail-of-simple-polygonal-path-is-simple
  by simp
  then have ih: simple-path (make-polygonal-path tl')
    using Suc.hyps * Suc.premis(1) Suc.premis(2) ab by force
  have simple-path ((linepath a b) +++ make-polygonal-path tl')
  proof-

```

```

let ?g1 = linepath a b
let ?g2 = make-polygonal-path tl'
let ?G1 = path-image ?g1
let ?G2 = path-image ?g2
have pathfinish ?g2 = last tl'
by (metis constant-linepath-is-not-loop-free ih last-conv-nth make-polygonal-path.simps(1)
polygon-pathfinish simple-path-def)
also have ... = vts ! (length vts - 2)
by (metis ab Suc.prem(1) Suc-1 constant-linepath-is-not-loop-free diff-Suc-1
diff-Suc-Suc ih impossible-Cons last.simps last-conv-nth length-Cons length-append-singleton
list.discI make-polygonal-path.simps(1) nle-le nth-append order-less-le simple-path-def)
finally have pathfinish-g2: pathfinish ?g2 = vts ! (length vts - 2) .

have pathfinish ?g1 = pathstart ?g2
by (metis ab constant-linepath-is-not-loop-free ih linepath-1' make-polygonal-path.simps(1)
pathfinish-def polygon-pathstart simple-path-def)
moreover have arc ?g1
by (metis Suc.prem(6) p arc-linepath constant-linepath-is-not-loop-free
not-loop-free-first-component simple-path-def)
moreover have arc ?g2
proof-
have pathstart ?g2 = b
using calculation(1) by auto
moreover have b = vts!1
by (metis ab One-nat-def Suc.prem(1) Suc.prem(2) Suc.prem(3)
Suc-le-eq length-append-singleton not-less-eq-eq nth-Cons-Suc nth-append numeral-3-eq-3)
moreover have last tl' ≠ vts!1
using loop-free-polygonal-path-vts-distinct Suc.prem
by (metis pre-ih ab append-Nil append-butlast-last-id butlast-conv-take but-
last-snoc calculation(2) constant-linepath-is-not-loop-free hd-conv-nth ih index-Cons
index-last list.collapse make-polygonal-path.simps(2) simple-path-def take0)
ultimately have pathfinish ?g2 ≠ b
using pathfinish-g2 ⟨pathfinish (make-polygonal-path tl') = last tl'⟩ by
presburger
thus ?thesis
using ⟨pathstart (make-polygonal-path tl') = b⟩ arc-simple-path ih by blast
qed
moreover have ?G1 ∩ ?G2 ⊆ {pathstart ?g2}
proof(rule subsetI)
let ?z = ((2::real)^(n-1) - 1)/(2^(n-1))
have g1: ?G1 = p{0..1/2}
proof-
have take 2 vts = [a, b]
by (smt (verit) * One-nat-def Suc.prem(1) Suc.prem(2) Suc-1 ab ap-
pend-Cons butlast-snoc drop0 drop-Suc-Cons length-append-singleton less-Suc-eq-le
not-less-eq-eq nth-butlast numeral-3-eq-3 plus-1-eq-Suc same-append-eq take-Suc-Cons
take-Suc-eq take-add take-all-iff)
then have ?g1 = make-polygonal-path (take 2 vts)
using make-polygonal-path.simps by presburger

```

```

moreover have  $1 < n$  using * by linarith
ultimately have  $?G1 = p\{0..(2^{2-1}) - 1)/(2^{2-1})\}$ 
  using vts-split-path-image
    by (metis * Suc.premis(2) Suc.premis(4) Suc-1 Suc-leD Suc-lessD
eval-nat-numeral(3) order.refl)
  thus ?thesis by force
qed
have  $g2: ?G2 = p\{1/2..?z\}$ 
proof-
  have  $tl' = take\ (n - 1)\ (drop\ 1\ vts)$ 
    using ab Suc.premis(1) Suc.premis(2) by simp
  moreover then have  $?g2 = make\ polygonal\ path\ (take\ (n - 1)\ (drop\ 1\ vts))$  by blast
  ultimately have  $?G2 = p\{(2^{1-1}) - 1)/(2^{1-1})..?z\}$ 
    using vts-sublist-path-image[of p vts ?g2 tl'  $n-1$   $1 - - n$   $((2::real)^1 - 1)/(2^1)$  ?z]
    by (metis * Suc.premis(1) Suc.premis(2) Suc.premis(4) Suc-eq-plus1
ab add-0 add-Suc-shift add-le-imp-le-diff diff-Suc-Suc diff-zero eval-nat-numeral(3)
length-Cons length-append less-Suc-eq-le list.size(3) order.refl)
  thus ?thesis by simp
qed
have  $1/2 \leq ?z$ 
  using * bounding-interval-helper1[of  $n-1$ ] Suc.premis
    by (smt (verit) One-nat-def diff-Suc-Suc less-diff-conv numeral-3-eq-3
one-le-power plus-1-eq-Suc power-one-right power-strict-increasing-iff real-shrink-le
add-2-eq-Suc diff-add-inverse less-trans-Suc numeral-eq-Suc pos2 self-le-power zero-less-diff)
  moreover have  $?z < 1$  by auto
  ultimately have  $z: 1/2 \leq ?z \wedge ?z < 1$  by blast

fix  $x$ 
assume  $x \in ?G1 \cap ?G2$ 
then obtain  $t1\ t2$  where  $t1t2: t1 \in \{0..1/2\} \wedge t2 \in \{1/2..?z\} \wedge p\ t1 = x \wedge p\ t2 = x$ 
  by (smt (verit, del-insts) g1 g2 Int-iff imageE path-image-def)
moreover have  $(t1 = t2) \vee (t1 = 0 \wedge t2 = 1) \vee (t1 = 1 \wedge t2 = 0)$ 
proof-
  have  $t1 \in \{0..1\} \wedge t2 \in \{0..1\}$ 
    by (meson t1t2  $z$  atLeastAtMost-iff dual-order.trans less-eq-real-def)
  thus ?thesis
  using Suc.premis(6) unfolding simple-path-def loop-free-def using t1t2
by presburger
qed
moreover have  $t1 = 1/2$  using calculation by force
ultimately have  $x = pathstart\ ?g2$ 
by (metis ab constant-linepath-is-not-loop-free dual-order.refl eq-divide-eq-numeral1(1)
ih joinpaths-def make-polygonal-path.simps(1) mult.commute p pathfinish-def pathfinish-linepath
polygon-pathstart simple-path-def zero-neq-numeral)
  thus  $x \in \{pathstart\ ?g2\}$  by simp
qed

```

```

ultimately show ?thesis using arc-join-eq ih by (metis arc-imp-simple-path)
qed
moreover have vts' = a # tl' using Suc.prem1 ab by argo
moreover have p' = (linepath a b) +++ make-polygonal-path tl'
proof -
  have Suc (length tl') = length vts' by (simp add: ab)
  then show ?thesis
    by (metis (no-types) * Cons-nth-drop-Suc Suc.prem1(1) Suc.prem1(2)
        Suc.prem1(5) Suc-lessD ab drop-0 length-append-singleton make-polygonal-path.simp1(4)
        not-less-eq numeral-3-eq-3)
  qed
ultimately have ?case by blast
}
ultimately show ?case using Suc.prem1 by linarith
qed

```

```

lemma take-i-is-loop-free:
  fixes vts :: (real^2) list
  assumes n = length vts
  assumes 2 ≤ i ∧ i ≤ n
  assumes vts' = take i vts
  assumes p = make-polygonal-path vts
  assumes p' = make-polygonal-path vts'
  assumes loop-free p
  shows loop-free p'
  using assms
proof(induct n-i arbitrary: vts' i p p')
  case 0
  moreover then have p = p' by auto
  ultimately show ?case by argo
next
  case (Suc x)

  let ?i' = i+1
  let ?q-vts = take (i+1) vts
  let ?q = make-polygonal-path ?q-vts

  have n-?i' = x using Suc.hyps(2) by linarith
  then have loop-free ?q using Suc.hyps Suc.prem1(2) Suc.prem1(4) Suc.prem1(6)
  assms(1) by auto
  moreover obtain z where ?q = make-polygonal-path (vts' @ [z])
  unfolding Suc.prem1(3)
  by (metis Suc.hyps(2) Suc-eq-plus1 assms(1) take-Suc-conv-app-nth zero-less-Suc
      zero-less-diff)
  ultimately show loop-free p'
  unfolding Suc.prem1 using one-append-simple-path unfolding simple-path-def
  by (metis One-nat-def Suc.prem1(2) Suc-1 add-diff-cancel-right' append-take-drop-id
      assms(1) diff-diff-cancel length-append length-append-singleton length-drop make-polygonal-path-gives-path
      not-less-eq-eq numeral-3-eq-3)

```

qed

lemma *sublist-is-loop-free*:

fixes $pts :: (\mathbb{R}^2) \text{ list}$
assumes $p = \text{make-polygonal-path } pts$
assumes $p' = \text{make-polygonal-path } pts'$
assumes *loop-free* p
assumes $m = \text{length } pts$
assumes $n = \text{length } pts'$
assumes *sublist* $pts' \ pts$
assumes $n \geq 2 \wedge m \geq 2$
shows *loop-free* p'

proof –

obtain $pre \ post$ **where** $pts = pre \ @ \ pts' \ @ \ post$ **using** *assms(6)* **unfolding**
sublist-def **by** *blast*
then have $pts' \ @ \ post = \text{drop } (\text{length } pre) \ pts$ **using** pts **by** *simp*
moreover have $pts' = \text{take } (\text{length } pts') \ (pts' \ @ \ post)$ **using** pts **by** *simp*
moreover have *loop-free* $(\text{make-polygonal-path } (pts' \ @ \ post))$
using *drop-i-is-loop-free* *assms* *calculation*
by (*smt* (*verit*, *del-insts*) *One-nat-def* *Suc-1* *Suc-leD* *diff-diff-cancel* *drop-all*
le-diff-iff' *length-append* *length-drop* *list.size(3)* *nat-le-linear* *not-numeral-le-zero*
numeral-3-eq-3 *trans-le-add1*)
ultimately show *?thesis*
using *take-i-is-loop-free* *assms*
by (*metis* *sublist-append-rightI* *sublist-length-le*)

qed

lemma *diff-points-path-image-set-property*:

fixes $a \ b :: \mathbb{R}^2$
assumes $a \neq b$
shows *path-image* $(\text{linepath } a \ b) \neq \{a, b\}$

proof –

have *not-a*: $(\text{linepath } a \ b) \ (1/2) \neq a$
by (*smt* (*verit*) *add-diff-cancel-left'* *assms* *divide-eq-0-iff* *linepath-def* *scaleR-cancel-left*
scaleR-collapse)
have *not-b*: $(\text{linepath } a \ b) \ (1/2) \neq b$
by (*smt* (*verit*, *ccfv-SIG*) *add-diff-cancel-right'* *assms* *divide-eq-1-iff* *linepath-def*
scaleR-cancel-left *scaleR-collapse*)
have $(\text{linepath } a \ b) \ (1/2) \in \text{path-image } (\text{linepath } a \ b)$
unfolding *path-image-def* **by** *simp*
then show *?thesis* **using** *not-a* *not-b* **by** *blast*

qed

lemma *polygonal-path-vertex-t*:

assumes $p = \text{make-polygonal-path } pts$
assumes $n = \text{length } pts$
assumes $n \geq 1$
assumes $0 \leq i \wedge i < n - 1$
assumes $x = (2^i - 1)/(2^i)$

```

shows vts!i = p x
using assms
proof(induct i arbitrary: p vts n x)
  case 0
  then show ?case
    by (metis bot-nat-0.extremum cancel-comm-monoid-add-class.diff-cancel diff-is-0-eq
div-0 less-nat-zero-code list.size(3) pathstart-def polygon-pathstart power-0)
  next
    case (Suc i)

  let ?vts' = tl vts
  let ?p' = make-polygonal-path ?vts'
  let ?x' =  $(2^i - 1)/(2^i)$ 

  have p x = ?p' ?x'
  proof-
    let ?a = vts!0
    let ?b = vts!1
    let ?l = linepath ?a ?b
    have  $n \geq 3$  using Suc.prems by linarith
    then have length ?vts' ≥ 2 by (simp add: Suc.prems(2))
    then have p = ?l +++ ?p'
      using Suc.prems make-polygonal-path.simps(4)[of ?a ?b ?vts!1 drop 2 ?vts']
      by (metis (no-types, opaque-lifting) Cons-nth-drop-Suc Suc-1 bot-nat-0.not-eq-extremum
diff-Suc-1 diff-is-0-eq drop-0 drop-Suc less-Suc-eq zero-less-diff)
    moreover have pathfinish ?l = pathstart ?p'
      by (metis One-nat-def  $\langle 2 \leq \text{length } (\text{tl } vts) \rangle$  length-greater-0-conv nth-tl order-less-le-trans pathfinish-linepath polygon-pathstart pos2)
    moreover have  $(\lambda x::\text{real}. 2 * x - 1) x = ?x'$ 
      using Suc.prems(5) Suc-eq-plus1 bounding-interval-helper1 diff-Suc-1 le-add2 summation-helper
      by presburger
    ultimately show ?thesis using joinpaths-tl-transform[of  $\lambda x. 2*x - 1$  ?l ?p' p
x]
      by (smt (verit, del-insts) divide-nonneg-nonneg half-bounded-equal two-realpow-ge-one)
    qed
    moreover have vts!(i+1) = ?vts!i using Suc.prems by (simp add: nth-tl)
    moreover have ?vts!i = ?p' ?x' using Suc.hyps Suc.prems by force
    ultimately show ?case by simp
  qed

```

```

lemma loop-free-split-int:
  assumes p = make-polygonal-path vts  $\wedge$  loop-free p
  assumes vts1 = take i vts
  assumes vts2 = drop (i-1) vts
  assumes c1 = make-polygonal-path vts1
  assumes c2 = make-polygonal-path vts2
  assumes  $n = \text{length } vts$ 
  assumes  $1 \leq i \wedge i < n$ 

```

shows $(\text{path-image } c1) \cap (\text{path-image } c2) \subseteq \{\text{pathstart } c1, \text{pathstart } c2\}$
 (is $?C1 \cap ?C2 \subseteq \{\text{pathstart } c1, \text{pathstart } c2\}$)
proof(*rule subsetI*)
let $?t = ((2::\text{real})^\wedge(i-1) - 1)/(2^\wedge(i-1))$

fix x
assume $x \in ?C1 \cap ?C2$
moreover have $c1c2: ?C1 = p\{0..?t\} \wedge ?C2 = p\{?t..1\}$
using *vts-split-path-image assms polygon-of-def by metis*
ultimately obtain $t1\ t2$ **where** $t1t2: t1 \in \{0..?t\} \wedge t2 \in \{?t..1\} \wedge p\ t1 = x$
 $\wedge p\ t2 = x$ **by** *auto*
moreover have $t1 \in \{0..1\} \wedge t2 \in \{0..1\}$ **using** *calculation by force*
moreover have $(t1 = t2) \vee (t1 = 0 \wedge t2 = 1)$
using *assms(1) calculation unfolding polygon-of-def polygon-def simple-path-def*
loop-free-def
by *fastforce*
ultimately have $x \in \{p\ ?t, p\ 0\}$ **by** *fastforce*
moreover have $p\ ?t = \text{pathstart } c2$
using *assms polygonal-path-vertex-t*
by (*smt (verit, cfv-SIG) Cons-nth-drop-Suc diff-less-mono le-eq-less-or-eq*
length-drop less-imp-diff-less less-trans-Suc less-zeroE linorder-neqE-nat list.size(3)
nth-Cons-0 numeral-1-eq-Suc-0 numerals(1) polygon-of-def polygon-pathstart)
moreover have $p\ 0 = \text{pathstart } c1$ **using** *assms*
by (*metis One-nat-def diff-is-0-eq diff-zero linorder-not-less nth-take path-*
start-def polygon-pathstart take-eq-Nil zero-less-Suc)
ultimately show $x \in \{\text{pathstart } c1, \text{pathstart } c2\}$ **by** *blast*
qed

lemma *loop-free-arc-split-int:*
assumes $p = \text{make-polygonal-path } vts \wedge \text{loop-free } p \wedge \text{arc } p$
assumes $vts1 = \text{take } i\ vts$
assumes $vts2 = \text{drop } (i-1)\ vts$
assumes $c1 = \text{make-polygonal-path } vts1$
assumes $c2 = \text{make-polygonal-path } vts2$
assumes $n = \text{length } vts$
assumes $1 \leq i \wedge i < n$
shows $(\text{path-image } c1) \cap (\text{path-image } c2) \subseteq \{\text{pathstart } c2\}$
 (is $?C1 \cap ?C2 \subseteq \{\text{pathstart } c2\}$)
proof(*rule subsetI*)
let $?t = ((2::\text{real})^\wedge(i-1) - 1)/(2^\wedge(i-1))$

fix x
assume $x \in ?C1 \cap ?C2$
moreover have $c1c2: ?C1 = p\{0..?t\} \wedge ?C2 = p\{?t..1\}$
using *vts-split-path-image assms polygon-of-def by metis*
ultimately obtain $t1\ t2$ **where** $t1t2: t1 \in \{0..?t\} \wedge t2 \in \{?t..1\} \wedge p\ t1 = x$
 $\wedge p\ t2 = x$ **by** *auto*
moreover have $t1 \in \{0..1\} \wedge t2 \in \{0..1\}$ **using** *calculation by force*
moreover have $(t1 = t2) \vee (t1 = 0 \wedge t2 = 1)$

using *assms(1) calculation unfolding polygon-of-def polygon-def simple-path-def loop-free-def*
by *fastforce*
moreover then have $t1 = t2$
using *assms(1) unfolding arc-def using calculation(1) inj-on-contrad by fastforce*
ultimately have $x \in \{p \ ?t\}$ **by** *fastforce*
moreover have $p \ ?t = \text{pathstart } c2$
using *assms polygonal-path-vertex-t*
by (*smt (verit, cfv-SIG) Cons-nth-drop-Suc diff-less-mono le-eq-less-or-eq length-drop less-imp-diff-less less-trans-Suc less-zeroE linorder-neqE-nat list.size(3) nth-Cons-0 numeral-1-eq-Suc-0 numerals(1) polygon-of-def polygon-pathstart*)
ultimately show $x \in \{\text{pathstart } c2\}$ **by** *fast*
qed

lemma *loop-free-append:*

assumes $p = \text{make-polygonal-path } vts$
assumes $p1 = \text{make-polygonal-path } vts1$
assumes $p2 = \text{make-polygonal-path } vts2$
assumes $vts = vts1 \ @ \ (tl \ vts2)$
assumes $\text{loop-free } p1 \ \wedge \ \text{loop-free } p2$
assumes $\text{path-image } p1 \ \cap \ \text{path-image } p2 \subseteq \{\text{pathstart } p1, \text{pathstart } p2\}$
assumes $\text{last } vts2 \neq \text{hd } vts1 \longrightarrow \text{path-image } p1 \ \cap \ \text{path-image } p2 \subseteq \{\text{pathstart } p2\}$
assumes $\text{last } vts1 = \text{hd } vts2$
assumes $\text{arc } p1 \ \wedge \ \text{arc } p2$
shows $\text{loop-free } p$
using *assms*
proof(*induct length vts1 arbitrary: p p1 p2 vts vts1 vts2 rule: less-induct*)
case *less*
have $1: \text{length } vts1 \geq 2$
using *less*
by (*metis Suc-1 arc-distinct-ends constant-linepath-is-not-loop-free diff-is-0-eq' make-polygonal-path.simps(1) not-less-eq-eq polygon-pathfinish polygon-pathstart*)
moreover have $\text{length } vts2 \geq 2$
using *less.prem*
by (*metis One-nat-def Suc-1 Suc-leI arc-distinct-ends diff-Suc-1 length-greater-0-conv make-polygonal-path.simps(1) nat-less-le pathfinish-linepath pathstart-linepath polygon-pathfinish polygon-pathstart*)
ultimately have $\text{length } vts \geq 3$ **using** *less assms(4) by auto*
{ assume $*$: $\text{length } vts1 = 2$
then obtain $a \ b$ **where** $vts1 = [a, b]$
by (*metis 1 Cons-nth-drop-Suc One-nat-def Suc-1 drop0 drop-eq-Nil lessI pos2*)
then have $p1: p1 = \text{linepath } a \ b$
using *less make-polygonal-path.simps(3) by metis*
have $p: p = p1 \ +++ \ p2$
using $p1$ *less*
by (*smt (verit) (vts1 = [a, b]) append-Cons assms(4) constant-linepath-is-not-loop-free last-ConsL last-ConsR list.exhaust-sel list.inject list.simps(3) make-polygonal-path.elims*)

```

self-append-conv2)
  have b: pathstart p2 ∈ path-image p1 ∩ path-image p2
    by (metis IntI less(3,4,6,9) constant-linepath-is-not-loop-free hd-conv-nth
last-conv-nth make-polygonal-path.simps(1) pathfinish-in-path-image pathstart-in-path-image
polygon-pathfinish polygon-pathstart)
    { assume pathstart p1 = pathfinish p2
      then have ?case using simple-path-join-loop-eq[of p2 p1] less.premis
        by (metis make-polygonal-path-gives-path p path-join-eq simple-path-def)
      } moreover
    { assume **: pathstart p1 ≠ pathfinish p2
      then have path-image p1 ∩ path-image p2 = {pathstart p2}
        using less.premis b
        by (metis constant-linepath-is-not-loop-free empty-subsetI hd-conv-nth in-
sert-subset last-conv-nth make-polygonal-path.simps(1) polygon-pathfinish polygon-pathstart
subset-antisym)
      then have ?case
        using arc-join-eq[of p1 p2]
        by (metis less(2,4,10) arc-imp-simple-path arc-join-eq-alt make-polygonal-path-gives-path
p path-join-path-ends simple-path-def)
      }
    ultimately have ?case by blast
  } moreover
  { assume *: length vts1 > 2
    then have len-p1: length vts1 ≥ 3 by linarith
    then obtain a b vts-tl where ab: vts = a # vts-tl ∧ b = hd vts-tl
      by (metis ‹3 ≤ length vts› length-0-conv list.collapse not-numeral-le-zero)
    have vts1-char: vts1 = (vts1 ! 0) # (vts1 ! 1) # (vts1 ! 2) # (drop 3 vts1)
      using len-p1
      by (metis 1 Cons-nth-drop-Suc One-nat-def Suc-1 drop0 length-greater-0-conv
linorder-not-less list.size(3) not-less-eq-eq not-numeral-le-zero numeral-3-eq-3)
    then have tail-vts1-char: tl vts1 = (vts1 ! 1) # (vts1 ! 2) # (drop 3 vts1)
      by (metis list.sel(3))

    let ?l = linepath a b
    let ?vts1-tl = tl vts1
    let ?p1-tl = make-polygonal-path ?vts1-tl
    let ?vts2-tl = tl vts2
    let ?p2-tl = make-polygonal-path ?vts2-tl
    let ?p-tl = make-polygonal-path vts-tl

    have p: p = ?l +++ ?p-tl
      unfolding less.premis(1)
      by (smt (verit, ccfv-SIG) Suc-le-length-iff ‹3 ≤ length vts› ab list.discI
list.sel(1) list.sel(3) make-polygonal-path.elims numeral-3-eq-3)
    have p1: p1 = ?l +++ ?p1-tl
      using ab unfolding less.premis(2)
      by (smt (verit, ccfv-SIG) * Nitpick.size-list-simp(2) One-nat-def Suc-1 Suc-le-eq
hd-append2 less.premis(4) list.sel(1) list.sel(3) make-polygonal-path.elims nat-less-le
tl-append2)
  }

```

have $p1\text{-img}$: $\text{path-image } ?l \cap \text{path-image } ?p1\text{-tl} = \{\text{pathstart } ?p1\text{-tl}\}$
by (*metis arc-join-eq-alt less.premis(2) less.premis(9) make-polygonal-path-gives-path p1 path-join-path-ends*)

have $vts\text{-tl} = ?vts1\text{-tl} @ (tl\ vts2)$
using $\text{less.premis}(4)$ *ab*
by (*metis * length-greater-0-conv list.sel(3) order.strict-trans pos2 tl-append2*)
moreover have $\text{loop-free } ?p1\text{-tl} \wedge \text{loop-free } p2$
using $\langle 3 \leq \text{length } vts1 \rangle \text{less.premis}(2) \text{less.premis}(5) \text{sublist-is-loop-free}$ **by** *fastforce*

moreover have $\text{path-image } ?p1\text{-tl} \cap \text{path-image } p2 \subseteq \{\text{pathstart } p2\}$
proof–
have $\text{path-image } ?p1\text{-tl} \subseteq \text{path-image } p1$
by (*metis (no-types, opaque-lifting) * Suc-1 Suc-lessD length-tl less.premis(2) list.collapse list.size(3) order.refl path-image-cons-union sup.bounded-iff zero-less-diff zero-order(3)*)
then have $\text{path-image } ?p1\text{-tl} \cap \text{path-image } p2 \subseteq \{\text{pathstart } p1, \text{pathstart } p2\}$
using *less* **by** *blast*
moreover have $\text{pathstart } p1 \notin \text{path-image } ?p1\text{-tl}$
proof(*rule ccontr*)
assume $\neg \text{pathstart } p1 \notin \text{path-image } ?p1\text{-tl}$
then have $\text{pathstart } p1 \in \text{path-image } ?p1\text{-tl}$ **by** *blast*
thus *False*
by (*metis (no-types, lifting) IntI arc-def arc-simple-path less(10) make-polygonal-path-gives-path p1 p1-img path-join-path-ends pathstart-in-path-image pathstart-join simple-path-joinE singletonD*)
qed
ultimately have $\text{path-image } ?p1\text{-tl} \cap \text{path-image } p2 \subseteq \{\text{pathstart } p2\}$ **by** *blast*
thus *?thesis* **by** *blast*
qed

moreover then have $\text{last } vts2 \neq \text{hd } ?vts1\text{-tl}$
 $\longrightarrow \text{path-image } ?p1\text{-tl} \cap \text{path-image } p2 \subseteq \{\text{pathstart } p2\}$ **by** *blast*
moreover have $\text{last } ?vts1\text{-tl} = \text{hd } vts2$
by (*metis * Suc-1 drop-Nil drop-Suc-Cons last-drop last-tl less.premis(8) list.collapse*)
moreover have $\text{arc } ?p1\text{-tl} \wedge \text{arc } p2$
by (*smt (verit, best) * Nitpick.size-list-simp(2) Suc-1 arc-imp-simple-path constant-linepath-is-not-loop-free diff-Suc-Suc diff-is-0-eq leD length-greater-0-conv length-tl less.premis(2) less.premis(5) less.premis(9) list.sel(3) make-polygonal-path.elims make-polygonal-path-gives-path order.strict-trans path-join-path-ends pos2 simple-path-joinE*)
ultimately have $ih1$: $\text{loop-free } ?p\text{-tl}$
using $\text{less.hyps}[\text{of } ?vts1\text{-tl } ?p\text{-tl } vts\text{-tl } ?p1\text{-tl } p2\ vts2]$ $*$ $\text{less.premis}(3)$ **by** *fastforce*

have $p\text{-tl}\text{-img}$: $\text{path-image } ?p\text{-tl} = \text{path-image } ?p1\text{-tl} \cup \text{path-image } p2$
by (*metis (no-types, lifting) * Suc-1 Suc-le-eq \langle 2 \leq \text{length } vts2 \rangle \langle \text{last } (tl\ vts1) = \text{hd } vts2 \rangle \langle vts\text{-tl} = tl\ vts1 @ tl\ vts2 \rangle \text{hd-conv-nth last-conv-nth length-greater-0-conv}*)

length-tl less.premis(3) less-diff-conv make-polygonal-path-image-append-alt order-less-le-trans path-image-join plus-1-eq-Suc polygon-pathfinish polygon-pathstart pos2)

have 1: *length* [a, b] < *length* vts1 **using** <3 ≤ *length* vts1> **by** *fastforce*
moreover have 2: *p* = *make-polygonal-path* vts **using** *less.premis(1)* **by** *auto*
moreover have 3: ?*l* = *make-polygonal-path* [a, b] **by** *simp*
moreover have 4: ?*p-tl* = *make-polygonal-path* vts-tl **using** *less* **by** *simp*
moreover have 5: vts = [a, b] @ tl vts-tl
using ab <3 ≤ *length* vts> *append-eq-Cons-conv* **by** *fastforce*
moreover have 6: *loop-free* ?*l* ∧ *loop-free* ?*p-tl*
proof–
have *sublist* [a, b] vts1
by (*metis* (*no-types*, *opaque-lifting*) 1 *Cons-nth-drop-Suc* *Suc-lessD* ab *append-Cons* *drop0* *length-Cons* *less.premis(4)* *list.sel(1)* *list.sel(3)* *list.size(3)* *sublist-take* *take0* *take-Suc-Cons*)
then have *loop-free* (*make-polygonal-path* [a, b])
using *sublist-is-loop-free* * *less.premis(2)* *less.premis(5)* **by** *fastforce*
then have *loop-free* ?*l* **using** *make-polygonal-path.simps(3)* **by** *simp*
thus ?*thesis* **using** *ih1* **by** *simp*
qed
moreover have 9: *last* [a, b] = *hd* vts-tl **by** (*simp* *add: ab*)
moreover have 10: *arc* ?*l* ∧ *arc* ?*p-tl*
proof–
have *pathstart* ?*p-tl* = b
by (*metis* 6 ab *constant-linepath-is-not-loop-free* *hd-conv-nth* *make-polygonal-path.simps(1)* *polygon-pathstart*)
moreover have *pathfinish* ?*p-tl* ≠ b
proof(*rule ccontr*)
assume ¬ *pathfinish* ?*p-tl* ≠ b
have *pathfinish* ?*p-tl* = *pathfinish* p2
by (*smt* (*verit*) 5 9 *Nil-tl* <2 ≤ *length* vts2> <¬ *pathfinish* (*make-polygonal-path* vts-tl) ≠ b> ab *arc-distinct-ends* *last-append* *last-conv-nth* *last-tl* *length-tl* *less.premis(3)* *less.premis(4)* *less.premis(9)* *list.size(3)* *not-numeral-le-zero* *polygon-pathfinish* *polygon-pathstart*)
moreover have b ∈ *path-image* p1
by (*metis* *list.size(3)* 1 *Cons-nth-drop-Suc* *Suc-lessD* *UnCI* ab *append-eq-conv-conj* *drop0* *hd-append2* *hd-conv-nth* *length-Cons* *less.premis(2)* *less.premis(4)* *list.distinct(1)* *list.sel(3)* *path-image-cons-union* *pathstart-in-path-image* *polygon-pathstart* *tl-append2*)
moreover have b ≠ *pathstart* p1
by (*metis* (*no-types*, *lifting*) 1 6 ab *constant-linepath-is-not-loop-free* *dual-order.strict-trans* *hd-append2* *hd-conv-nth* *length-greater-0-conv* *less.premis(2)* *less.premis(4)* *list.sel(1)* *list.size(3)* *polygon-pathstart*)
moreover have b ≠ *pathfinish* p2
by (*metis* (*no-types*, *lifting*) *Int-insert-right-if1* *arc-distinct-ends* *calculation(2)* *calculation(3)* *insert-absorb* *insert-iff* *insert-not-empty* *less.premis(6)* *less.premis(9)* *pathfinish-in-path-image* *subset-iff*)
ultimately show *False*
using <¬ *pathfinish* (*make-polygonal-path* vts-tl) ≠ b> **by** *fastforce*
qed

```

ultimately have pathstart ?p-tl ≠ pathfinish ?p-tl by simp
then have arc ?p-tl
  using ih1 arc-def loop-free-cases make-polygonal-path-gives-path by metis
moreover have arc ?l by (metis 6 arc-linepath constant-linepath-is-not-loop-free)
ultimately show ?thesis by blast
qed
moreover have 7: path-image ?l ∩ path-image ?p-tl ⊆ {pathstart ?l, pathstart
?p-tl}
proof-
  have path-image ?l ⊆ path-image p1
    by (metis Un-iff ⟨loop-free (make-polygonal-path (tl vts1)) ∧ loop-free
p2⟩ ⟨vts-tl = tl vts1 @ tl vts2⟩ ab constant-linepath-is-not-loop-free hd-append2
hd-conv-nth make-polygonal-path.simps(1) p1 path-image-join pathfinish-linepath
polygon-pathstart subsetI)
  then have path-image ?l ∩ path-image p2 ⊆ {pathstart p1, pathstart p2}
    using less.premis(6) by auto
  moreover have pathstart p2 ∉ path-image ?l
    by (smt (verit, ccfv-threshold) 10 Int-insert-left-if1 ⟨arc (make-polygonal-path
(tl vts1)) ∧ arc p2⟩ ⟨last (tl vts1) = hd vts2⟩ ⟨loop-free (make-polygonal-path (tl
vts1)) ∧ loop-free p2⟩ arc-def arc-distinct-ends arc-join-eq-alt constant-linepath-is-not-loop-free
hd-conv-nth insert-absorb last-conv-nth less.premis(3) less.premis(9) make-polygonal-path.simps(1)
p1 path-join-eq pathfinish-in-path-image polygon-pathfinish polygon-pathstart single-
ton-insert-inj-eq)
  ultimately have path-image ?l ∩ path-image ?p-tl ⊆ {pathstart p1, pathstart
?p1-tl}
    using p1-img p-tl-img by blast
  moreover have pathstart ?p1-tl = pathstart ?p-tl
  by (metis 2 less.premis(2) make-polygonal-path-gives-path p p1 path-join-path-ends)
  moreover have pathstart p1 = pathstart ?l by (simp add: p1)
  ultimately show ?thesis by argo
qed
moreover have 8: last vts-tl ≠ hd [a, b]
  → path-image ?l ∩ path-image ?p-tl ⊆ {pathstart ?p-tl}
proof clarify
  fix x
  assume a1: last vts-tl ≠ hd [a, b]
  assume a2: x ∈ path-image ?l
  assume a3: x ∈ path-image ?p-tl

  have hd vts1 ≠ last vts2
    using less.premis
  by (metis a1 ⟨vts-tl = tl vts1 @ tl vts2⟩ ab arc-distinct-ends constant-linepath-is-not-loop-free
hd-append2 last-appendR last-tl length-tl list.sel(1) list.size(3) make-polygonal-path.simps(1)
polygon-pathfinish polygon-pathstart)
  then have p1-p2-int: path-image p1 ∩ path-image p2 ⊆ {pathstart p2}
    using less.premis by argo

  have x ≠ pathstart ?l
  proof(rule ccontr)

```

```

assume **:  $\neg x \neq \text{pathstart } ?l$ 
have  $\text{pathstart } ?l \notin \text{path-image } ?p1\text{-tl}$ 
by (metis Int-iff arc-distinct-ends arc-join-eq-alt empty-iff insertE less.premis(2)
less.premis(9) make-polygonal-path-gives-path p1 path-join-path-ends pathstart-in-path-image)
then have  $\text{pathstart } ?l \in \text{path-image } p2$  using p1-img p-tl-img ** a3 by
blast
then have  $\text{pathstart } ?l \in \text{path-image } p1 \cap \text{path-image } p2$ 
by (metis IntI p1 pathstart-in-path-image pathstart-join)
moreover have  $\text{pathstart } ?l \neq \text{pathstart } p2$ 
by (metis arc-distinct-ends constant-linepath-is-not-loop-free hd-conv-nth
last-conv-nth less.premis(2) less.premis(3) less.premis(5) less.premis(8) less.premis(9)
make-polygonal-path.simps(1) p1 pathstart-join polygon-pathfinish polygon-pathstart)
ultimately show False using p1-p2-int by blast
qed
moreover have  $x = \text{pathstart } ?l \vee x = \text{pathstart } ?p\text{-tl}$  using 7 a2 a3 by
blast
ultimately show  $x = \text{pathstart } ?p\text{-tl}$  by fast
qed
ultimately have ?case using less.hyps[of [a, b] p vts ?l ?p-tl vts-tl] by blast
}
ultimately show ?case using less 1 by linarith
qed

```

lemma *sublist-path-image-subset*:

```

assumes sublist vts1 vts2
assumes length vts1  $\geq 1$ 
shows  $\text{path-image } (\text{make-polygonal-path } vts1) \subseteq \text{path-image } (\text{make-polygonal-path } vts2)$ 
proof –
let ?p1 =  $\text{make-polygonal-path } vts1$ 
let ?p2 =  $\text{make-polygonal-path } vts2$ 
let ?m =  $\text{length } vts1$ 
let ?n =  $\text{length } vts2$ 
have n-geq-m: ?n  $\geq$  ?m by (simp add: assms(1) sublist-length-le)

have ?thesis if *:  $\text{length } vts1 = 1$ 
proof –
have  $\text{path-image } ?p1 = \{vts1!0\}$ 
by (metis Cons-nth-drop-Suc One-nat-def closed-segment-idem drop0 drop-eq-Nil
le-numeral-extra(4) make-polygonal-path.simps(2) path-image-linepath that zero-less-one)
moreover have  $vts1!0 \in \text{set } vts2$ 
by (metis assms(1) less-numeral-extra(1) nth-mem set-mono-sublist subsetD
that)
ultimately show ?thesis
using vertices-on-path-image by force
qed
moreover have ?thesis if *:  $\text{length } vts1 \geq 2$ 
proof –
obtain pre post where sublist:  $vts2 = \text{pre } @ vts1 @ \text{post}$ 

```

```

    using assms(1) unfolding sublist-def by blast
  let ?i = length pre
  let ?j = length vts1
  let ?k = ?i + ?j
  let ?x1 = (2?i - 1) / 2(?i)::real
  let ?x2 = (2(?k-1) - 1) / (2(?k-1))::real
  let ?x = (2(?i-1) - 1) / 2(?i-1)::real
  have path-image ?p1 = ?p2 ‘ {?x1..?x2} if **: length post ≥ 1
    using sublist * ** vts-sublist-path-image[of ?p2 vts2 ?p1 vts1 ?j ?i ?n ?m ?k
    ?x1 ?x2]
    by fastforce
  moreover have path-image ?p1 = ?p2 ‘ {?x1..1} if **: length post = 0
  proof-
    have sublist: vts2 = pre @ vts1 using ** sublist by blast
    moreover have vts1 = drop ?i vts2 using sublist * by simp
    moreover have 1 ≤ ?i + 1 ∧ ?i + 1 < length vts2 using sublist * ** by
  simp
    ultimately show ?thesis
      using vts-split-path-image[of ?p2 vts2 - - ?p1 vts1 ?i + 1 ?n ?x1] add-diff-cancel-right’
      by metis
    qed
    moreover have ?p2 ‘ {?x1..?x2} ⊆ path-image ?p2 ∧ ?p2 ‘ {?x1..1} ⊆
  path-image ?p2
    proof-
      have {?x1..?x2} ⊆ {0..1} ∧ {?x1..1} ⊆ {0..1} by simp
      thus ?thesis unfolding path-image-def by blast
    qed
    ultimately show ?thesis by (metis less-one linorder-not-le)
  qed
  ultimately show ?thesis using assms by linarith
  qed

```

```

lemma integral-on-edge-subset-integral-on-path:
  assumes p = make-polygonal-path vts and
    (i::int) ∈ {0.. $((length\ vts) - 1)$ } and
    x = vts!i and
    y = vts!(i+1)
  shows {v. integral-vec v ∧ v ∈ path-image (linepath x y)}
    ⊆ {v. integral-vec v ∧ v ∈ path-image p}
  using assms edge-subset-path-image by blast

```

```

lemma sublist-pair-integral-subset-integral-on-path:
  assumes p = make-polygonal-path vts and
    sublist [x, y] vts
  shows {v. integral-vec v ∧ v ∈ path-image (linepath x y)}
    ⊆ {v. integral-vec v ∧ v ∈ path-image p}
  using assms integral-on-edge-subset-integral-on-path
  proof-
    obtain pre post where vts: vts = pre @ [x, y] @ post using assms(2) sublist-def

```

by *blast*
let $?i = \text{length } \text{pre}$
have $x = \text{vts}! ?i$ **using** vts **by** *simp*
moreover **have** $y = \text{vts}!(?i + 1)$
by (*metis vts add.right-neutral append-Cons nth-Cons-Suc nth-append-length*
nth-append-length-plus plus-1-eq-Suc)
moreover **have** $?i \in \{0..<((\text{length } \text{vts}) - 1)\}$ **using** vts **by** *force*
ultimately **show** $?thesis$ **using** $\text{assms}(1)$ *integral-on-edge-subset-integral-on-path*
by *auto*
qed

lemma *sublist-integral-subset-integral-on-path*:

assumes $\text{length } \text{ell} \geq 2$
assumes $p = \text{make-polygonal-path } \text{vts}$ **and**
sublist ell vts
shows $\{v. \text{integral-vec } v \wedge v \in \text{path-image } (\text{make-polygonal-path } \text{ell})\}$
 $\subseteq \{v. \text{integral-vec } v \wedge v \in \text{path-image } p\}$
proof –
obtain $\text{pre } \text{post}$ **where** $\text{vts} : \text{vts} = \text{pre} @ \text{ell} @ \text{post}$ **using** $\text{assms}(3)$ *sublist-def*
by *blast*
then **have** $\text{len-vts} : \text{length } \text{vts} \geq 2$
using $\text{assms}(1)$
by *auto*
let $?i = \text{length } \text{pre}$
have $v \in \text{path-image } p$ **if** $*$: $v \in \text{path-image } (\text{make-polygonal-path } \text{ell})$ **for** v
proof –
have $\exists j :: \text{nat}. v \in \text{path-image } (\text{linepath } (\text{ell} ! j) (\text{ell} ! (j+1))) \wedge j+1 < \text{length } \text{ell}$
using $*$ *polygonal-path-image-linepath-union assms(1)*
by (*meson less-diff-conv make-polygonal-path-image-property*)
then **obtain** j **where** $v\text{-in} : v \in \text{path-image } (\text{linepath } (\text{ell} ! j) (\text{ell} ! (j+1)))$
 $j+1 < \text{length } \text{ell}$
by *auto*
then **have** $\text{ell-at} : \text{ell} ! j = \text{vts} ! (j + \text{length } \text{pre}) \wedge \text{ell} ! (j+1) = \text{vts} ! (j + 1 + \text{length } \text{pre})$
using vts
by (*simp add: nth-append*)
then **have** $v\text{-in}2 : v \in \text{path-image } (\text{linepath } (\text{vts} ! (j + \text{length } \text{pre})) (\text{vts} ! (j + \text{length } \text{pre} + 1)))$
using $v\text{-in}(1)$ **by** *simp*
have $j + 1 + \text{length } \text{pre} < \text{length } \text{vts}$
using $\text{ell-at } v\text{-in}(2)$ vts **by** *auto*
then **have** $j\text{-plus} : j + \text{length } \text{pre} < \text{length } \text{vts} - 1$
by *auto*
then **show** $?thesis$ **using** $v\text{-in}2$ *linepaths-subset-make-polygonal-path-image[OF len-vts j-plus]* $\text{assms}(1)$
 $\text{assms}(2)$ **by** *auto*
qed
then **show** $?thesis$ **by** *blast*

qed

13 Reversing Polygonal Path Vertex List

lemma *rev-vts-path-image*:
shows $\text{path-image } (\text{make-polygonal-path } (\text{rev } \text{vts})) = \text{path-image } (\text{make-polygonal-path } \text{vts})$

proof –

- { **assume** $\text{length } \text{vts} \leq 1$
- then have** *?thesis*
- by** (*smt* (*verit*, *best*) *One-nat-def Suc-length-conv le-SucE le-zero-eq length-0-conv rev.simps(1) rev-singleton-conv*)
- } **moreover**
- { **fix** x
- assume** $*$: $x \in \text{path-image } (\text{make-polygonal-path } (\text{rev } \text{vts})) \wedge \text{length } \text{vts} \geq 2$
- then obtain** k **where** $k\text{-prop}$: $k < \text{length } (\text{rev } \text{vts}) - 1 \wedge x \in \text{path-image } (\text{linepath } (\text{rev } \text{vts} ! k) (\text{rev } \text{vts} ! (k + 1)))$
- using** *make-polygonal-path-image-property*[*of rev vts*] **by** *auto*
- have** $p1$: $\text{rev } \text{vts} ! k = \text{vts} ! (\text{length } \text{vts} - k - 1)$
- using** *rev-nth*
- by** (*metis* *Suc-lessD* $\langle k < \text{length } (\text{rev } \text{vts}) - 1 \wedge x \in \text{path-image } (\text{linepath } (\text{rev } \text{vts} ! k) (\text{rev } \text{vts} ! (k + 1))) \rangle$ *add.commute diff-diff-left length-rev less-diff-conv plus-1-eq-Suc*)
- have** $p2$: $\text{rev } \text{vts} ! (k + 1) = \text{vts} ! (\text{length } \text{vts} - k - 2)$
- using** *rev-nth*[*of k+1 vts*] $k\text{-prop}$
- by** *force*
- then have** $x \in \text{path-image } (\text{linepath } (\text{vts} ! (\text{length } \text{vts} - k - 1)) (\text{vts} ! (\text{length } \text{vts} - k - 2)))$
- using** $k\text{-prop}$ $p1$ $p2$ **by** *auto*
- then have** $x \in \text{path-image } (\text{linepath } (\text{vts} ! (\text{length } \text{vts} - k - 2)) (\text{vts} ! (\text{length } \text{vts} - k - 1)))$
- using** *reversepath-linepath path-image-reversepath*
- by** *metis*
- then have** $x \in \text{path-image } (\text{make-polygonal-path } \text{vts})$
- using** *linepaths-subset-make-polygonal-path-image* $*$ $k\text{-prop}$
- by** (*smt* (*verit*, *best*) *Nat.diff-add-assoc add.commute add-diff-cancel-left' diff-le-self length-rev less-Suc-eq less-diff-conv linorder-not-less nat-1-add-1 nat-neq-iff plus-1-eq-Suc subsetD*)
- } **moreover**
- { **fix** x
- assume** $*$: $x \in \text{path-image } (\text{make-polygonal-path } \text{vts}) \wedge \text{length } \text{vts} \geq 2$
- then obtain** k **where** $k\text{-prop}$: $k < \text{length } \text{vts} - 1 \wedge x \in \text{path-image } (\text{linepath } (\text{vts} ! k) (\text{vts} ! (k + 1)))$
- using** *make-polygonal-path-image-property*[*of vts*] **by** *auto*
- have** $p1$: $\text{vts} ! k = (\text{rev } \text{vts}) ! (\text{length } \text{vts} - k - 1)$
- using** *rev-nth* $k\text{-prop}$
- by** (*metis* *Suc-eq-plus1 Suc-lessD diff-diff-left length-rev less-diff-conv rev-rev-ident*)
- have** $p2$: $\text{vts} ! (k + 1) = (\text{rev } \text{vts}) ! (\text{length } \text{vts} - k - 2)$
- using** *rev-nth*[*of k+1*]

```

    by (smt (verit) Suc-eq-plus1 add-2-eq-Suc' diff-diff-left k-prop length-rev
less-diff-conv rev-rev-ident)
    then have x ∈ path-image (linepath (rev vts ! (length vts - k - 2)) (rev vts !
(length vts - k - 1)))
    using reversepath-linepath path-image-reversepath
    by (metis k-prop p1)
    then have x ∈ path-image (make-polygonal-path (rev vts))
    using linepaths-subset-make-polygonal-path-image k-prop *
    by (smt (verit, best) Suc-1 Suc-diff-Suc Suc-eq-plus1 Suc-le-eq Suc-lessD
bot-nat-0.not-eq-extremum diff-commute diff-diff-left diff-less length-rev less-numeral-extra(1)
subsetD zero-less-diff)
  }
  ultimately show ?thesis by force
qed

```

```

lemma rev-vts-is-loop-free:
  assumes p = make-polygonal-path vts
  assumes loop-free p
  shows loop-free (make-polygonal-path (rev vts))
  using assms
proof(induct length vts arbitrary: p vts)
  case 0
  then show ?case by simp
next
  case (Suc n)
  then have Suc n ≥ 2
  by (metis One-nat-def Suc-length-conv constant-linepath-is-not-loop-free le-SucE
le-add1 le-numeral-Suc length-greater-0-conv list.size(3) make-polygonal-path.simps(2)
numeral-One plus-1-eq-Suc pred-numeral.simps(2) semiring-norm(26))
  moreover
  { assume *: Suc n = 2
    then obtain a b where ab: p = linepath a b
    using Suc.premis make-polygonal-path.simps(3)
    by (metis (no-types, opaque-lifting) Cons-nth-drop-Suc One-nat-def Suc.hyps(2)
Suc-1 diff-Suc-1 drop-0 drop-Suc length-0-conv length-tl zero-less-Suc)
    moreover then have a ≠ b using Suc.premis(2) constant-linepath-is-not-loop-free
  by blast
    ultimately have loop-free (linepath b a) by (simp add: linepath-loop-free)
    moreover have make-polygonal-path (rev vts) = linepath b a
    by (smt (z3) * Cons-nth-drop-Suc One-nat-def Suc.hyps(2) Suc.premis(1)
Suc-1 Suc-diff-Suc ab butlast-snoc diff-Suc-1 drop0 hd-conv-nth hd-rev last-conv-nth
length-butlast length-rev lessI linepath-1' make-polygonal-path.simps(3) nth-append-length
pathstart-def pathstart-linepath pos2 rev.simps(2) rev-is-Nil-conv rev-take take-eq-Nil)
    ultimately have ?case by simp
  } moreover
  { assume *: Suc n > 2
    let ?vts' = butlast vts
    let ?p' = make-polygonal-path ?vts'
    let ?vts'-rev = rev ?vts'

```

```

let ?p'-rev = make-polygonal-path ?vts'-rev

let ?vts-rev = rev vts
let ?p-rev = make-polygonal-path ?vts-rev

obtain y z where yz: y = last ?vts'  $\wedge$  z = last vts by blast
let ?l = linepath y z
let ?l-rev = linepath z y
have loop-free ?p'
by (metis * Suc.hyps(2) Suc.prem(1) Suc.prem(2) butlast-conv-take diff-Suc-1
le-add2 less-Suc-eq-le plus-1-eq-Suc take-i-is-loop-free)
then have loop-free-p'-rev: loop-free ?p'-rev using Suc.hyps by force
moreover have rev vts = z # ?vts'-rev
by (metis Suc.hyps(2) yz append-butlast-last-id length-0-conv nat.distinct(1)
rev-eq-Cons-iff rev-rev-ident)
moreover have y = hd ?vts'-rev using yz by (simp add: hd-rev)
ultimately have p-rev: ?p-rev = ?l-rev +++ ?p'-rev
by (smt (verit, best) constant-linepath-is-not-loop-free list.sel(1) make-polygonal-path.elims
make-polygonal-path.simps(4))

have [y, z] = drop (n-1) vts
using yz Suc.hyps(2)
by (metis (no-types, opaque-lifting) * Cons-nth-drop-Suc Suc-1 Suc-diff-Suc
Suc-lessD Suc-n-not-le-n append-butlast-last-id append-eq-conv-conj diff-Suc-1 last-conv-nth
length-0-conv length-butlast less-nat-zero-code linorder-not-le nth-take)
then have ?l = make-polygonal-path (drop (n-1) vts)
using make-polygonal-path.simps by metis
moreover have ?p' = make-polygonal-path (take n vts)
using Suc.hyps(2) by (metis butlast-conv-take diff-Suc-1)
ultimately have path-image ?l  $\cap$  path-image ?p'  $\subseteq$  {pathstart ?l, pathstart
?p'}
using loop-free-split-int
by (smt (verit, ccfv-SIG) Int-commute Suc.hyps(2) Suc.prem(1) Suc.prem(2)
Suc-1 Suc-le-mono  $\langle 2 \leq$  Suc n  $\rangle$  insert-commute lessI)
moreover have path-image ?l = path-image ?l-rev by auto
moreover have path-image ?p' = path-image ?p'-rev
using * Suc.hyps(2) rev-vts-path-image by force
moreover have pathstart ?l = pathfinish ?l-rev by simp
moreover have pathstart ?p' = pathfinish ?p'-rev
by (metis Nil-is-rev-conv last.simps last-conv-nth last-rev list.distinct(1)
list.exhaust-sel make-polygonal-path.simps(1) make-polygonal-path.simps(2) nth-Cons-0
polygon-pathfinish polygon-pathstart)
ultimately have path-image-int:
path-image ?l-rev  $\cap$  path-image ?p'-rev  $\subseteq$  {pathfinish ?l-rev, pathfinish
?p'-rev}
by argo

have 1: pathfinish ?l-rev = pathstart ?p'-rev
by (metis make-polygonal-path-gives-path p-rev path-join-path-ends)

```

```

{ assume pathfinish ?p'-rev = pathstart ?l-rev
  then have ?case using simple-path-join-loop 1 p-rev path-image-int
    by (smt (verit, del-insts) Suc.hyps(2) Suc.prem(1) Suc.prem(2) Suc-1
      ‹linepath y z = make-polygonal-path (drop (n - 1) vts)› ‹loop-free (make-polygonal-path
        (rev (butlast vts)))› constant-linepath-is-not-loop-free diff-Suc-Suc drop-i-is-loop-free
        dual-order.eq-iff insert-commute linepath-loop-free make-polygonal-path-gives-path
        path-linepath pathfinish-linepath pathstart-linepath simple-path-cases simple-path-def)
    } moreover
  { assume pathfinish ?p'-rev ≠ pathstart ?l-rev
    then have pathstart p ≠ pathfinish p
      by (metis Suc.prem(1) ‹loop-free (make-polygonal-path (butlast vts))› ‹path-
        start (make-polygonal-path (butlast vts)) = pathfinish (make-polygonal-path (rev
          (butlast vts)))› butlast-conv-take constant-linepath-is-not-loop-free last-conv-nth less-nat-zero-code
          make-polygonal-path.simps(1) nat-neq-iff nth-take pathstart-linepath polygon-pathfinish
          polygon-pathstart take-eq-Nil yz)
    } then have arc p
      by (metis Suc.prem(1) Suc.prem(2) arc-def loop-free-cases make-polygonal-path-gives-path)
    then have path-image ?l-rev ∩ path-image ?p'-rev ⊆ {pathstart ?p'-rev}
      using loop-free-arc-split-int
      by (metis 1 Int-commute Suc.hyps(2) Suc.prem(1) Suc.prem(2) ‹2 ≤ Suc
        n› ‹linepath y z = make-polygonal-path (drop (n - 1) vts)› ‹make-polygonal-path
          (butlast vts) = make-polygonal-path (take n vts)› ‹path-image (linepath y z) =
            path-image (linepath z y)› ‹path-image (make-polygonal-path (butlast vts)) = path-image
              (make-polygonal-path (rev (butlast vts)))› ‹pathstart (linepath y z) = pathfinish
                (linepath z y)› le-numeral-Suc lessI numerals(1) pred-numeral-simps(2) semiring-norm(26))
    } moreover have arc ?l-rev
      by (metis Suc.hyps(2) Suc.prem(1) Suc.prem(2) Suc-1 ‹[y, z] = drop (n -
        1) vts› arc-linepath constant-linepath-is-not-loop-free diff-Suc-Suc drop-i-is-loop-free
        dual-order.refl make-polygonal-path.simps(3))
    moreover have arc ?p'-rev
      proof–
      have ?p'-rev 0 = last (butlast vts) by (metis 1 pathfinish-linepath pathstart-def
        yz)
      moreover have ?p'-rev 1 = hd (butlast vts)
      by (metis ‹loop-free (make-polygonal-path (butlast vts))› ‹pathstart (make-polygonal-path
        (butlast vts)) = pathfinish (make-polygonal-path (rev (butlast vts)))› constant-linepath-is-not-loop-free
        hd-conv-nth make-polygonal-path.simps(1) pathfinish-def polygon-pathstart)
      moreover have last (butlast vts) ≠ hd (butlast vts) using Suc.prem
      by (metis (no-types, lifting) * Suc.hyps(2) Suc-1 diff-is-0-eq index-Cons
        index-last leD length-butlast less-diff-conv less-imp-le-nat list.collapse list.size(3)
        loop-free-polygonal-path-vts-distinct not-one-le-zero plus-1-eq-Suc)
      ultimately have ?p'-rev 0 ≠ ?p'-rev 1 by simp
      thus ?thesis using loop-free-p'-rev
      by (metis arc-def loop-free-cases make-polygonal-path-gives-path pathfin-
        ish-def pathstart-def)
    } qed
  ultimately have ?case
    using arc-join-eq[OF 1] arc-imp-simple-path p-rev simple-path-def by auto
  }

```

```

    ultimately have ?case by blast
  }
  ultimately show ?case by linarith
qed

```

```

lemma rev-vts-is-polygon:
  assumes polygon-of p vts
  shows polygon (make-polygonal-path (rev vts))
  using rev-vts-is-loop-free assms
  unfolding polygon-of-def polygon-def simple-path-def
  using make-polygonal-path-gives-path
  by (metis One-nat-def closed-path-def UNIV-def length-greater-0-conv polygon-pathfinish
  polygon-pathstart polygonal-path-def rangeI rev.simps(1) rev-nth rev-rev-ident)

```

```

end
theory Linepath-Collinearity
  imports Polygon-Lemmas

```

```

begin

```

14 Collinearity Properties

```

lemma points-on-linepath-collinear:
  assumes exists-c: ( $\exists c. a - b = c *_R u$ )
  assumes x-in-linepath:  $x \in \text{path-image (linepath a b)}$ 
  shows ( $\exists c. x - a = c *_R u$ ) ( $\exists c. b - x = c *_R u$ )
proof -
  obtain k :: real where k-prop:  $0 \leq k \wedge k \leq 1 \wedge x = (1 - k) *_R a + k *_R b$ 
    using x-in-linepath unfolding linepath-def path-image-def by fastforce
  then have  $x = a - k *_R a + k *_R b$ 
    by (simp add: eq-diff-eq)
  then have  $x - a = -k *_R a + k *_R b$ 
    by auto
  then have  $x - a = -k *_R (a - b)$ 
    by (simp add: scaleR-right-diff-distrib)
  obtain c where c-prop:  $a - b = c *_R u$  using exists-c by blast
  show ( $\exists c. x - a = c *_R u$ ) using  $x - a = -k *_R (a - b)$ 
    by (metis scaleR-scaleR)
  then show ( $\exists c. b - x = c *_R u$ )
    using exists-c
    by (metis (no-types, opaque-lifting) add-diff-eq diff-add-cancel minus-diff-eq
  scaleR-left-distrib)
qed

```

```

lemma three-points-collinear-property:
  fixes a b:: real^2
  assumes exists-c1: ( $\exists c. a - x1 = c *_R u$ )
  assumes exists-c2: ( $\exists c. a - x2 = c *_R u$ )
  shows  $\exists c. x1 - x2 = c *_R u$ 

```

```

proof –
  obtain c1 where c1-prop:  $a - x1 = c1 *_R u$ 
    using exists-c1 by auto
  obtain c2 where c2-prop:  $a - x2 = c2 *_R u$ 
    using exists-c2 by auto
  then have  $a - x2 - (a - x1) = c2 *_R u - c1 *_R u$ 
    using c1-prop c2-prop by simp
  then have  $a - x2 - (a - x1) = (c2 - c1) *_R u$ 
    by (simp add: scaleR-left-diff-distrib)
  then show ?thesis
    by auto
qed

lemma in-path-image-imp-collinear:
  fixes a b::  $\text{real}^2$ 
  assumes  $k \in \text{path-image } (\text{linepath } a \ b)$ 
  shows collinear  $\{a, b, k\}$ 
proof –
  obtain w where w-prop:  $w \in \{0..1\} \wedge k = (1 - w) *_R a + w *_R b$ 
    using assms unfolding path-image-def linepath-def by fast
  have collinear  $\{0, a-b, (1 - w) *_R a + (w-1) *_R b\}$ 
    using collinear
  by (smt (verit) collinear-lemma diff-minus-eq-add scaleR-minus-left scaleR-right-diff-distrib)
  then have collinear  $\{0, a - b, k - b\}$ 
    using w-prop
  by (metis (no-types, lifting) add.commute add-diff-cancel-left collinear-lemma
scaleR-collapse scaleR-right-diff-distrib)
  then show ?thesis using assms collinear-alt collinear-3 [of a b k]
    by auto
qed

lemma two-linepath-colinearity-property:
  fixes a b c d::  $\text{real}^2$ 
  assumes  $y \neq z \wedge \{y, z\} \subseteq (\text{path-image } (\text{linepath } a \ b)) \cap (\text{path-image } (\text{linepath } c \ d))$ 
  shows collinear  $\{a, b, c, d\}$ 
proof –
  have collinear  $\{a, b, y, z\}$ 
    using in-path-image-imp-collinear assms
  by (metis (no-types, lifting) Int-closed-segment collinear-4-3 inf.boundedE inf-idem
insert-absorb2 insert-subset path-image-linepath pathstart-in-path-image pathstart-linepath)
  moreover have collinear  $\{c, d, y, z\}$ 
    using in-path-image-imp-collinear assms
  by (metis (no-types, lifting) Int-closed-segment collinear-4-3 inf.boundedE inf-idem
insert-absorb2 insert-subset path-image-linepath pathstart-in-path-image pathstart-linepath)
  ultimately show ?thesis
    using assms collinear-3-eq-affine-dependent collinear-4-3 insert-absorb2 insert-commute
    by (smt (z3) collinear-3-trans)

```

qed

lemma *polygon-vts-not-collinear*:

assumes *polygon-of* p vts

shows \neg *collinear* (*set* vts)

proof –

have *len-vts*: $length\ vts \geq 3$

using *polygon-at-least-3-vertices* *assms* **unfolding** *polygon-of-def*

using *card-length* *dual-order.trans* **by** *blast*

have *compact-and-connected*: $compact\ (path\ image\ p) \wedge connected\ (path\ image\ p)$

using *inside-outside-polygon* *assms* **unfolding** *polygon-of-def*

using *compact-simple-path-image* *connected-simple-path-image* *polygon-def*

by *auto*

have *nonempty-path-image*: $path\ image\ p \neq \{\}$

using *assms* **unfolding** *polygon-of-def*

using *vertices-on-path-image* **by** *simp*

have *collinear-imp*: $collinear\ (set\ vts) \implies (collinear\ (path\ image\ p))$

proof –

assume *collinear* (*set* vts)

then obtain u where *u-prop*: $\forall x \in set\ vts. \forall y \in set\ vts. \exists c. x - y = c *_{\mathbb{R}} u$

unfolding *collinear-def* **by** *blast*

then have $\exists c. x - y = c *_{\mathbb{R}} u$ **if** *xy-in-pathimage*: $y \in path\ image\ p \wedge x \in path\ image\ p$ **for** $x\ y$

proof –

obtain $k1$ where *k1-prop*: $k1 < length\ vts - 1 \wedge x \in path\ image\ (linepath\ (vts\ !\ k1)\ (vts\ !\ (k1 + 1)))$

using *make-polygonal-path-image-property* *xy-in-pathimage* *len-vts*

by (*metis* *One-nat-def* *Suc-1* *Suc-leD* *assms* *numeral-3-eq-3* *polygon-of-def*)

then have $\exists c. (vts\ !\ k1) - (vts\ !\ (k1 + 1)) = c *_{\mathbb{R}} u$

by (*meson* *add-lessD1* *in-set-conv-nth* *less-diff-conv* *u-prop*)

obtain $k2$ where *k2-prop*: $k2 < length\ vts - 1 \wedge y \in path\ image\ (linepath\ (vts\ !\ k2)\ (vts\ !\ (k2 + 1)))$

using *make-polygonal-path-image-property* *xy-in-pathimage* *len-vts*

by (*metis* *One-nat-def* *Suc-1* *Suc-leD* *assms* *numeral-3-eq-3* *polygon-of-def*)

have $\exists c. vts\ !\ (k2 + 1) - (vts\ !\ k1) = c *_{\mathbb{R}} u$

using *u-prop* *k1-prop* *k2-prop*

by (*meson* *add-lessD1* *less-diff-conv* *nth-mem*)

have *k2-vts-prop*: $\exists c. vts\ !\ (k2 + 1) - (vts\ !\ k2) = c *_{\mathbb{R}} u$

using *u-prop* *k2-prop* **by** *fastforce*

have *ex-c-k2*: $\exists c. vts\ !\ (k2 + 1) - y = c *_{\mathbb{R}} u$

using *points-on-linepath-collinear*[*of* $vts\ !\ (k2 + 1)\ vts\ !\ k2\ u\ y$] *k2-prop* *k2-vts-prop*

by (*meson* *add-lessD1* *points-on-linepath-collinear*(2) *less-diff-conv* *nth-mem* *u-prop*)

have *k1-vts-prop*: $\exists c. vts\ !\ (k1 + 1) - (vts\ !\ k1) = c *_{\mathbb{R}} u$

using *u-prop* *k1-prop* **by** *fastforce*

have *ex-c-k1-y*: $\exists c. vts\ !\ (k1 + 1) - y = c *_{\mathbb{R}} u$

using *points-on-linepath-collinear*[*of* $vts\ !\ (k1 + 1)\ vts\ !\ k1\ u\ y$] *k1-prop*

k1-vts-prop
by (*meson* $\langle \exists c. vts ! (k2 + 1) - vts ! k1 = c *_R u \rangle \langle \exists c. vts ! k1 - vts ! (k1 + 1) = c *_R u \rangle$ *three-points-collinear-property ex-c-k2*)
have *ex-c-k1-x*: $\exists c. vts ! (k1 + 1) - x = c *_R u$
using *points-on-linepath-collinear*[*of vts ! (k1 + 1) vts ! k1 u x*] *k1-prop*
k1-vts-prop
by (*meson* *add-lessD1 points-on-linepath-collinear(2) less-diff-conv nth-mem u-prop*)
show *?thesis*
using *ex-c-k1-y ex-c-k1-y three-points-collinear-property ex-c-k1-x* **by** *blast*
qed
then show (*collinear (path-image p)*) **unfolding** *collinear-def* **by** *auto*
qed
{ **assume** *: *collinear (set vts)*
then obtain *a b::real^2* **where** *im-closed: path-image p = closed-segment a b*
using *collinear-imp compact-convex-collinear-segment-alt*[*of path-image p*]
compact-and-connected nonempty-path-image
by *blast*
have *inside (closed-segment a b) = {}*
by (*simp add: inside-convex*)
then have *path-inside p = {}*
unfolding *path-inside-def* **using** *im-closed* **by** *auto*
then have *False*
using *inside-outside-polygon assms* **unfolding** *polygon-of-def inside-outside-def*
by *blast*
}
then show *?thesis* **by** *blast*
qed

lemma *not-collinear-with-subset*:
assumes *collinear A*
assumes \neg *collinear (A \cup {x})*
assumes *card A > 2*
assumes *a \in A*
shows \neg *collinear ((A - {a}) \cup {x})*
proof–
obtain *u v* **where** *uv: u \in A \wedge v \in A \wedge u \neq v \wedge u \neq a \wedge v \neq a*
proof–
have *card (A - {a}) \geq 2* **using** *assms* **by** *auto*
then obtain *u B* **where** *u \in (A - {a}) \wedge B = (A - {a}) - {u}*
by (*metis bot-nat-0.extremum-unique card.empty ex-in-conv zero-neq-numeral*)
moreover then obtain *v* **where** *v \in B*
by (*metis Diff-iff One-nat-def Suc-1 assms(3) assms(4) card.empty card.insert equalsOI finite.intros(1) finite-insert insert-Diff insert-commute less-irrefl*)
ultimately show *?thesis* **using** *that* **by** *blast*
qed
then have *x \notin affine hull {u, v}*
using *assms*
by (*smt (verit, ccfv-threshold) Un-commute Un-upper1 collinear-affine-hull-collinear*

hull-insert hull-mono insert-absorb insert-is-Un insert-subset
moreover have $u \in A - \{a\} \wedge v \in A - \{a\}$ **using** *uv* **by** *blast*
ultimately show *?thesis*
by (*metis UnCI collinear-3-imp-in-affine-hull collinear-triples insert-absorb singletonD uv*)
qed

lemma *vec-diff-scale-collinear*:

fixes $a\ b\ c :: \text{real}^2$
assumes $b - a = m *_R (c - a)$
shows *collinear* $\{a, b, c\}$
proof –
{ **assume** $m = 0$
then have $b = a$ **using** *assms* **by** *simp*
then have *collinear* $\{a, b, c\}$ **by** *auto*
} **moreover**
{ **assume** *m-nz*: $m \neq 0$
then have *c-eq*: $c = (1/m) *_R (b - a) + a$ **using** *assms* **by** *simp*
then have $c - b = (1/m - 1) *_R (b - a)$ **using** *m-nz* **by** (*simp add: scaleR-left.diff*)
then obtain m' **where** $c - b = m' *_R (b - a)$ **by** *fast*

then have $c - b \in \text{span}(\{b - a\})$ **by** (*simp add: span-breakdown-eq*)
moreover from *this* **have** $b - c \in \text{span}(\{b - a\})$ **using** *span-0 span-add-eq2*
by *fastforce*
moreover have $c - a \in \text{span}(\{b - a\})$ **using** *assms* **by** (*simp add: span-breakdown-eq c-eq*)
moreover from *this* **have** $a - c \in \text{span}(\{b - a\})$ **using** *span-0 span-add-eq2*
by *fastforce*
moreover have $b - a \in \text{span}(\{b - a\})$ **by** (*simp add: span-base*)
moreover from *this* **have** $a - b \in \text{span}(\{b - a\})$ **using** *span-0 span-add-eq2*
by *fastforce*
moreover have $\forall v \in \{a, b, c\}. v - v \in \text{span}(\{b - a\})$ **by** (*simp add: span-0*)
ultimately have $\forall v \in \{a, b, c\}. \forall w \in \{a, b, c\}. v - w \in \text{span}(\{b - a\})$ **by**
blast
then have $\forall v \in \{a, b, c\}. \forall w \in \{a, b, c\}. \exists k. v - w = k *_R (b - a)$
by (*simp add: span-breakdown-eq*)
then have *collinear* $\{a, b, c\}$ **using** *collinear-def* **by** *blast*
} **ultimately show** *?thesis* **using** *assms* **by** *auto*
qed

15 Linepath Properties

lemma *good-linepath-comm*: $\text{good-linepath } a\ b\ vts \implies \text{good-linepath } b\ a\ vts$
unfolding *good-linepath-def*
by (*metis (no-types, opaque-lifting) insert-commute path-image-linepath segment-convex-hull*)

lemma *finite-set-linepaths*:
assumes *polygon*: *polygon* *p*
assumes *polygonal-path*: $p = \text{make-polygonal-path } vts$
shows *finite* $\{(a, b). (a, b) \in \text{set } vts \times \text{set } vts\}$
proof –
have *finite* (*set vts*)
using *polygonal-path* **by** *auto*
then have *finite* (*set vts* \times *set vts*)
by *blast*
then show *?thesis*
by *auto*
qed

lemma *linepaths-intersect-once-or-collinear*:
fixes *a b c d* :: real^2
assumes *path-image* (*linepath* *a b*) \cap *path-image* (*linepath* *c d*) $\neq \{\}$
shows *collinear* $\{a, b, c, d\} \vee (\exists x. \text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d) = \{x\})$
proof *safe*
assume $\neg (\exists x. \text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d) = \{x\})$
then obtain *x y* **where** $x \neq y \wedge \{x, y\} \subseteq \text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d)$
using *assms* **by** *blast*
then show *collinear* $\{a, b, c, d\}$ **using** *two-linepath-collinearity-property* **by** *meson*
qed

lemma *linepaths-intersect-once-or-collinear-alt*:
fixes *a b c d* :: real^2
assumes *path-image* (*linepath* *a b*) \cap *path-image* (*linepath* *c d*) $\neq \{\}$
shows *collinear* $\{a, b, c, d\} \vee \text{card } (\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d)) = 1$
proof –
have $\text{card } (\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d)) = 1$
 $\longleftrightarrow (\exists x. \text{path-image } (\text{linepath } a \ b) \cap \text{path-image } (\text{linepath } c \ d) = \{x\})$
using *is-singleton-altdef* *is-singleton-def* **by** *blast*
thus *?thesis* **using** *linepaths-intersect-once-or-collinear* *assms* **by** *presburger*
qed

lemma *path-image-linepath-union*:
fixes *a b* :: *'a::euclidean-space*
assumes *d* \in *path-image* (*linepath* *a b*)
shows *path-image* (*linepath* *a b*) = *path-image* (*linepath* *a d*) \cup *path-image* (*linepath* *d b*)
proof –
have *path-image* (*linepath* *a b*) = *closed-segment* *a b* **using** *path-image-linepath*
by *simp*
also then have $\dots = \text{closed-segment } a \ d \cup \text{closed-segment } d \ b$
using *Un-closed-segment* *assms* **by** *blast*

```

also have ... = path-image (linepath a d) ∪ path-image (linepath d b)
  using path-image-linepath by simp
ultimately show ?thesis by order
qed

lemma path-image-linepath-split:
  assumes i < (length vts) - 1
  assumes x ∈ path-image (linepath (vts!i) (vts!(i+1)))
  assumes x-notin: x ∉ set vts
  shows path-image (make-polygonal-path vts) = path-image (make-polygonal-path
    ((take (i+1) vts) @ [x] @ (drop (i+1) vts)))
  using assms
proof(induct length vts arbitrary: vts i x)
  case 0
  then show ?case by linarith
next
  case (Suc n)
  let ?vts' = (take (i+1) vts) @ [x] @ (drop (i+1) vts)
  let ?p = make-polygonal-path vts
  let ?p' = make-polygonal-path ?vts'
  have Suc n ≥ 2 using Suc by linarith
  then obtain v1 v2 vts-tail where vts-is: vts = v1 # v2 # vts-tail
  by (metis Suc(2) Cons-nth-drop-Suc One-nat-def Suc-1 Suc-le-eq drop0 zero-less-Suc)

  { assume *: i = 0
    then have vts'-is: ?vts' = [v1, x, v2] @ vts-tail
      using vts-is by simp
    then have x-in: x ∈ path-image (linepath v1 v2)
      using * Suc.prem1 vts-is by simp
    { assume *: vts-tail = []
      then have p-is: path-image ?p = path-image (linepath v1 v2)
        using vts-is make-polygonal-path.simps(3)[of v1 v2]
        by simp
      have path-image ?p' = path-image (linepath v1 x) ∪ path-image (linepath x
v2)
        using vts'-is * make-polygonal-path.simps(4)[of v1 x v2 []]
        using make-polygonal-path.simps(3)[of x v2]
        by (metis append.right-neutral list.discI nth-Cons-0 path-image-cons-union)
      then have ?case
        using p-is path-image-linepath-union[of x v1 v2] assms(3) vts-is x-in by
blast
    } moreover
    { assume *: vts-tail ≠ []
      then have path-image ?p = path-image (linepath v1 v2) ∪ path-image
(make-polygonal-path (v2 # vts-tail))
        using path-image-cons-union vts-is by (metis list.discI nth-Cons-0)
      moreover have path-image (linepath v1 x) ∪ path-image (linepath x v2) =
path-image (linepath v1 v2)
        using path-image-linepath-union x-in by blast
    }
  }

```

```

    ultimately have ?case
      by (metis (no-types, lifting) append-Cons append-Nil inf-sup-aci(6) list.discI
nth-Cons-0 path-image-cons-union vts'-is)
    }
    ultimately have ?case by blast
  } moreover
  { assume * : i > 0
    then have Suc n > 2 using Suc by linarith

    let ?vts-tl = tl vts
    let ?vts-tl' = (take i ?vts-tl) @ [x] @ (drop i ?vts-tl)
    let ?p-tl = make-polygonal-path ?vts-tl
    let ?p-tl' = make-polygonal-path ?vts-tl'

    have ?vts-tl!(i-1) = vts!i ∧ ?vts-tl!i = vts!(i+1) using Suc * by (simp add:
vts-is)
    moreover then have x ∈ path-image (linepath (?vts-tl!(i-1)) (?vts-tl!i))
      using Suc by presburger
    ultimately have path-image ?p-tl = path-image ?p-tl'
      using Suc
      by (smt (verit) * One-nat-def Suc-leI diff-Suc-1 le-add-diff-inverse2 length-tl
less-diff-conv list.sel(3) list.set-intros(2) vts-is)
    moreover have path-image ?p = path-image (linepath v1 v2) ∪ path-image
?p-tl
      using path-image-cons-union vts-is by auto
    ultimately have ?case
      by (smt (verit, ccfv-threshold) Nil-is-append-conv Suc-eq-plus1 ⟨i = 0 ⇒
path-image (make-polygonal-path vts) = path-image (make-polygonal-path (take (i
+ 1) vts @ [x] @ drop (i + 1) vts))⟩ append-Cons append-same-eq append-take-drop-id
drop-Suc hd-append2 hd-conv-nth list.sel(1) list.sel(3) path-image-cons-union take-eq-Nil
vts-is)
    }
    ultimately show ?case by linarith
  qed

```

```

lemma linepath-split-is-loop-free:
  assumes d ∈ path-image (linepath a b)
  assumes d ∉ {a, b}
  shows loop-free (make-polygonal-path [a, d, b]) (is loop-free ?p)
proof -
  let ?l1 = linepath a d
  let ?l2 = linepath d b
  have path-image ?l1 ∩ path-image ?l2 = {d} using Int-closed-segment assms(1)
by auto
  moreover have arc ?l1 ∧ arc ?l2 using assms(2) by fastforce
  ultimately show ?thesis
    by (metis arc-imp-simple-path arc-join-eq-alt make-polygonal-path.simps(3)
make-polygonal-path.simps(4) pathfinish-linepath pathstart-linepath simple-path-def)
  qed

```

```

lemma loop-free-linepath-split-is-loop-free:
  assumes  $p = \text{make-polygonal-path } vts$ 
  assumes loop-free  $p$ 
  assumes  $n = \text{length } vts$ 
  assumes  $i < n - 1$ 
  assumes  $x \in \text{path-image } (\text{linepath } (vts!i) (vts!(i+1))) \wedge x \notin \text{set } vts$ 
  assumes  $vts' = (\text{take } (i+1) vts) @ [x] @ (\text{drop } (i+1) vts)$ 
  assumes  $p' = \text{make-polygonal-path } vts'$ 
  shows loop-free  $p' \wedge \text{path-image } p' = \text{path-image } p$ 
  using assms
proof (induct  $i$  arbitrary:  $p$   $vts$   $p'$   $vts' n$ )
  case 0
  let  $?vts\text{-tl} = \text{tl } vts$ 
  let  $?p\text{-tl} = \text{make-polygonal-path } ?vts\text{-tl}$ 
  let  $?vts'\text{-tl} = \text{tl } vts'$ 
  let  $?p'\text{-tl} = \text{make-polygonal-path } ?vts'\text{-tl}$ 
  let  $?a = vts!0$ 
  let  $?b = vts!1$ 
  let  $?l = \text{linepath } ?a ?b$ 
  let  $?l' = \text{make-polygonal-path } [?a, x, ?b]$ 

  have  $vts!i: vts' = [?a, x] @ ?vts\text{-tl}$ 
  using 0
  by (metis (no-types, lifting) Suc-eq-plus1 append-Cons append-eq-append-conv2
append-self-conv bot-nat-0.not-eq-extremum diff-is-0-eq drop0 drop-Suc list.collapse
nth-Cons-0 take-Suc take-all-iff take-eq-Nil)

  have  $x \notin \{?a, ?b\}$ 
  by (metis 0(3-5) One-nat-def Suc-eq-plus1 bot-nat-0.not-eq-extremum diff-is-0-eq
insert-iff less-diff-conv nth-mem singletonD take-Suc-eq take-all-iff)
  then have lf-l': loop-free  $?l'$  using linepath-split-is-loop-free[of  $x$   $?a$   $?b$ ] 0 by
simp

  { assume  $\text{length } ?vts\text{-tl} = 1$ 
    then have  $vts' = [?a, x, ?b]$ 
    by (metis Cons-nth-drop-Suc One-nat-def append-eq-Cons-conv drop0 drop-eq-Nil
le-numeral-extra(4) nth-tl  $vts'$  zero-less-one)
    then have  $?case$  using linepath-split-is-loop-free path-image-linepath-split
    by (metis 0.prems(1) 0.prems(3) 0.prems(4) 0.prems(5) 0.prems(6) 0.prems(7)
lf-l')
  } moreover
  { assume  $*$ :  $\text{length } ?vts\text{-tl} \geq 2$ 
    then have  $p: p = ?l +++ ?p\text{-tl}$ 
    using make-polygonal-path.simps(4)[of  $?a$   $?b$ ]
    by (metis (no-types, opaque-lifting) 0(1) 0(3) 0(4) Cons-nth-drop-Suc
One-nat-def Suc-1 Suc-le-eq diff-is-0-eq drop-0 drop-Suc length-tl less-nat-zero-code
nat-le-linear nth-tl)
  }

```

```

have loop-free ?p-tl
  using tail-of-loop-free-polygonal-path-is-loop-free 0 *
  by (metis list.exhaust-sel list.sel(2))
moreover have l-l': path-image ?l = path-image ?l'
  using path-image-linepath-split 0
  by (metis One-nat-def Suc-eq-plus1 list.discI make-polygonal-path.simps(3)
nth-Cons-0 path-image-cons-union path-image-linepath-union)
moreover have path-image ?l' ∩ path-image ?p-tl ⊆ {?a, ?b}
  by (metis (mono-tags, opaque-lifting) p l-l' 0.prem(1) 0.prem(2) make-polygonal-path-gives-path
path-join-path-ends pathfinish-linepath pathstart-linepath simple-path-def simple-path-joinE)
moreover have arc p → path-image ?l' ∩ path-image ?p-tl ⊆ {?b}
  using p l-l'
  by (metis arc-def arc-join-eq make-polygonal-path-gives-path path-join-eq
path-linepath pathfinish-linepath)
moreover have arc p ↔ hd [?a, x, ?b] ≠ last (tl vts)
  by (metis * 0.prem(1) 0.prem(2) arc-def arc-simple-path last-conv-nth last-tl
list.sel(1) list.sel(2) list.size(3) loop-free-cases make-polygonal-path-gives-path not-numeral-le-zero
polygon-pathfinish polygon-pathstart)
moreover have vts' = [?a, x, ?b] @ tl ?vts-tl
  by (metis drop-Suc 0.prem(3) 0.prem(4) One-nat-def append-Cons ap-
pend-Nil append-take-drop-id length-tl nth-tl take-Suc-conv-app-nth take-eq-Nil vts')
moreover have last [?a, x, ?b] = hd ?vts-tl
  by (metis 0.prem(3) 0.prem(4) One-nat-def hd-conv-nth last.simps length-greater-0-conv
length-tl list.discI nth-tl)
moreover have pathfinish ?l = pathstart ?p-tl
  by (metis (no-types) 0.prem(1) make-polygonal-path.simps(3) make-polygonal-path-gives-path
p path-join-eq)
moreover have ∧v va vb vs. pathfinish (linepath v va) = pathstart (make-polygonal-path
(va # vb # vs))

  by (metis (no-types) make-polygonal-path.simps(3) make-polygonal-path.simps(4)
make-polygonal-path-gives-path path-join-eq)
ultimately have loop-free p'
  using loop-free-append[of p' vts' ?l' [?a, x, ?b] ?p-tl ?vts-tl]
  by (metis (no-types) 0.prem(1) 0.prem(2) 0.prem(7) arc-simple-path lf-l'
make-polygonal-path.simps(3) make-polygonal-path.simps(4) make-polygonal-path-gives-path
p pathfinish-join pathstart-linepath simple-path-def simple-path-joinE)
  then have ?case
    using 0(1) 0(3) 0(4) 0(5) 0(6) 0(7) path-image-linepath-split by blast
  }
ultimately show ?case
  by (metis 0(3,4) One-nat-def Suc-lessI length-tl less-eq-Suc-le nat-1-add-1
plus-1-eq-Suc)
next
case (Suc i)
let ?vts-tl = tl vts
let ?p-tl = make-polygonal-path ?vts-tl
let ?vts'-tl = tl vts'
let ?p'-tl = make-polygonal-path ?vts'-tl

```

```

let ?a = vts!0
let ?b = vts!1
let ?l = linepath ?a ?b

have ?vts-tl!i = vts!(Suc i) ∧ ?vts-tl!(i+1) = vts!((Suc i) + 1)
  by (metis Suc.premis(3) Suc.premis(4) add-Suc-right add-Suc-shift diff-is-0-eq
linorder-not-le list.exhaust-sel list.size(3) not-less-zero nth-Cons-Suc)
moreover have set ?vts-tl ⊆ set vts
  by (metis list.sel(2) list.set-sel(2) subsetI)
ultimately have x ∈ path-image (linepath (?vts-tl!i) (?vts-tl!(i+1))) ∧ x ∉ set
?vts-tl
  using Suc.premis(5) by auto
moreover have vts'-tl: ?vts'-tl = (take (i+1) ?vts-tl) @ [x] @ (drop (i+1)
?vts-tl)
  by (metis Suc.premis(3) Suc.premis(4) Suc.premis(6) Suc-eq-plus1 drop-Suc leD
length-tl take-all-iff take-eq-Nil take-tl tl-append2 zero-eq-add-iff-both-eq-0 zero-neq-one)
moreover have loop-free ?p-tl
  using tail-of-loop-free-polygonal-path-is-loop-free Suc.premis
by (metis Nitpick.size-list-simp(2) Suc-1 Suc-leI Suc-neq-Zero diff-0-eq-0 diff-Suc-1
less-one linorder-neqE-nat list.collapse not-less-zero)
ultimately have ih: loop-free ?p'-tl ∧ path-image ?p'-tl = path-image ?p-tl
  using Suc.premis Suc.hyps[of ?p-tl ?vts-tl - ?vts'-tl ?p'-tl] by simp

have p: p = ?l +++ ?p-tl
proof -
  have f1: ∀ vs. (hd (tl vs)::(real, 2) vec) = vs ! 1 ∨ [] = vs ∨ [] = tl vs
    by (metis (no-types) One-nat-def hd-conv-nth list.collapse nth-Cons-Suc)
  have [] ≠ tl vts ∧ vts ≠ [] ∧ tl vts ≠ [hd (tl vts)]
    by (metis Suc.premis(1) Suc.premis(2) ‹loop-free (make-polygonal-path (tl vts))›
constant-linepath-is-not-loop-free make-polygonal-path.simps(1) make-polygonal-path.simps(2))
  then have p = make-polygonal-path [hd vts, vts ! 1] +++ make-polygonal-path
(tl vts) ∧ vts ≠ []
    using f1 by (metis (full-types) Suc.premis(1) list.collapse make-polygonal-path.simps(3)
make-polygonal-path.simps(4))
  then show ?thesis
    by (simp add: hd-conv-nth)
qed

have length vts' ≥ 3 using Suc.premis by force
moreover have ab: ?a = vts!0 ∧ ?b = vts!1
  using Suc.premis
  by (smt (verit, ccfv-SIG) One-nat-def Suc-eq-plus1 add-Suc-right append-Cons
drop0 drop-Suc length-tl less-nat-zero-code list.exhaust-sel list.size(3) nat-diff-split
nth-Cons-0 nth-Cons-Suc take-Suc zero-less-Suc)
ultimately have p': p' = ?l +++ ?p'-tl
  using Suc.premis(7) make-polygonal-path.simps(4)[of ?a ?b]
  by (metis (no-types, opaque-lifting) Cons-nth-drop-Suc One-nat-def Suc-leD
Suc-le-eq drop0 drop-Suc numeral-3-eq-3)

```

```

have nonarc: path-image ?l  $\cap$  path-image ?p-tl  $\subseteq$  {?a, ?b}
  using simple-path-join-loop-eq Suc.prem3
  by (smt (verit, ccfv-threshold) p One-nat-def length-tl less-zeroE make-polygonal-path-gives-path
nth-tl order.strict-iff-not order-le-less-trans path-join-eq path-linepath pathfinish-linepath
pathstart-linepath polygon-pathstart simple-path-def simple-path-joinE take-Nil take-all-iff)
  have arc: arc p  $\longrightarrow$  path-image ?l  $\cap$  path-image ?p-tl  $\subseteq$  {?b}
  using arc-join-eq
  by (metis Suc.prem3(1) p make-polygonal-path-gives-path path-join-eq path-linepath
pathfinish-linepath)

{ assume arc p
  moreover then have path-image ?l  $\cap$  path-image ?p'-tl  $\subseteq$  {?b} using arc ih
by presburger
  moreover have pathfinish ?l = pathstart ?p'-tl
  by (metis Suc.prem3(7) make-polygonal-path-gives-path p' path-join-path-ends)
  ultimately have ?case using p' arc-join-eq[of ?l ?p'-tl]
  by (smt (verit, ccfv-SIG) Nil-is-append-conv Suc.prem3(3) Suc.prem3(4)
Suc-eq-plus1 vts'-tl arc-simple-path drop-eq-Nil ih last-appendR last-conv-nth last-drop
leD length-tl make-polygonal-path-gives-path p path-image-join path-join-eq path-linepath
pathfinish-linepath polygon-pathfinish simple-path-def simple-path-joinE take-all-iff
take-eq-Nil)
} moreover
{ assume  $\neg$  arc p
  then have pathstart ?l = pathfinish ?p'-tl  $\wedge$  pathfinish ?l = pathstart ?p'-tl
  by (smt (verit, del-insts) Nil-is-append-conv Nil-tl One-nat-def Suc.prem3(2)
Suc.prem3(3) Suc.prem3(4) Suc-eq-plus1 vts'-tl ab arc-def drop-eq-Nil last-appendR
last-conv-nth last-drop leD length-tl list.collapse loop-free-cases make-polygonal-path-gives-path
nth-Cons-Suc p path-join-eq path-linepath pathfinish-join pathfinish-linepath path-
start-join polygon-pathfinish polygon-pathstart take-all-iff take-eq-Nil)
  then have ?case using simple-path-join-loop-eq[of ?l ?p'-tl] p' nonarc
  by (smt (verit, ccfv-threshold) One-nat-def Suc.prem3(2) Suc.prem3(3) Suc.prem3(4)
arc-def constant-linepath-is-not-loop-free dual-order.strict-trans ih leD length-tl loop-free-cases
make-polygonal-path-gives-path not-loop-free-first-component nth-tl p path-image-join
path-linepath pathfinish-linepath pathstart-linepath polygon-pathstart simple-path-def
simple-path-join-loop-eq take-all-iff take-eq-Nil zero-less-Suc)
}
ultimately show ?case by argo
qed

```

```

lemma polygon-linepath-split-is-polygon:
  assumes polygon-of p vts
  assumes  $i < (\text{length } vts) - 1$ 
  assumes  $a = vts!i \wedge b = vts!(i+1)$ 
  assumes  $x \in \text{path-image } (\text{linepath } a \ b) \wedge x \notin \text{set } vts$ 
  assumes  $vts' = (\text{take } (i+1) \ vts) @ [x] @ (\text{drop } (i+1) \ vts)$ 
  shows polygon (make-polygonal-path vts')
proof -
  let ?p' = make-polygonal-path vts'

```

```

have path ?p' using assms make-polygonal-path-gives-path by presburger
moreover have loop-free ?p' using assms loop-free-linepath-split-is-loop-free
  by (metis polygon-def polygon-of-def simple-path-def)
moreover have closed-path ?p'
proof -
  have hd vts' = hd vts
  using assms
  by (metis hd-append2 hd-take le-diff-conv linorder-not-less take-all-iff take-eq-Nil2
trans-less-add2 zero-less-one)
  moreover have last vts' = last vts
  using assms linordered-semidom-class.add-diff-inverse by auto
  ultimately show ?thesis
  by (metis closed-path-def ⟨path ?p'⟩ append-butlast-last-id append-eq-conv-conj
append-is-Nil-conv assms(1) assms(5) have-wraparound-vertex hd-conv-nth length-butlast
not-Cons-self nth-append-length polygon-of-def polygon-pathfinish polygon-pathstart)
  qed
  ultimately show ?thesis unfolding polygon-def polygonal-path-def simple-path-def
assms(5) by blast
qed

```

16 Measure of linepaths

lemma *linepath-is-negligible-vertical*:

```

fixes a b :: real^2
assumes a$1 = b$1
defines p ≡ linepath a b
shows negligible (path-image p)
proof -
  have p-t: ∀ t ∈ {0..1}. (p t)$1 = a$1
  using linepath-in-path p-def segment-vertical assms by blast

```

```

let ?x = a$1
let ?e1 = (vector [1, 0])::real^2

```

```

have (1::real) ∈ Basis by simp
then have axis 1 (1::real) ∈ (⋃ i. ⋃ u ∈ (Basis::(real set)). {axis i u}) by blast
moreover have ?e1 = axis 1 (1::real)
  unfolding axis-def vector-def by auto
ultimately have e1-basis: ?e1 ∈ (Basis::((real^2) set)) by simp
then have negligible {v. v · ?e1 = ?x} (is negligible ?S)
  using negligible-standard-hyperplane by auto
moreover have ∀ t ∈ {0..1}. (p t) · ?e1 = ?x

```

```

proof clarify
  fix t :: real
  assume t: t ∈ {0..1}
  have (p t) · ?e1 = (p t)$1
  by (smt (verit, best) e1-basis cart-eq-inner-axis vec-nth-Basis vector-2(1))
  also have ... = ?x using p-t t by blast
  finally show (p t) · ?e1 = ?x .

```

qed
 moreover from *this* have *path-image* $p \subseteq ?S$ **unfolding** *path-image-def* by *blast*
 ultimately show *?thesis* **using** *negligible-subset* by *blast*
 qed

lemma *linepath-is-negligible-non-vertical*:

fixes $a\ \$1\ b\ ::\ real^{\wedge}2$
 assumes $a\ \$1\ <\ b\ \1
 defines $p \equiv linepath\ a\ b$
 shows *negligible* (*path-image* p)
proof –
 let $?A = (vector\ [vector\ [1,\ b\ \$1 - a\ \$1],\ vector\ [0,\ b\ \$2 - a\ \$2]])::(real^{\wedge}2^{\wedge}2)$
 let $?f1 = \lambda v::real^{\wedge}2. (?A * v)$
 let $?id = \lambda v::real^{\wedge}2. v$
 let $?f-a = \lambda v::real^{\wedge}2. a$
 let $?f2 = \lambda v. ?id\ v + ?f-a\ v$
 let $?f = ?f2 \circ ?f1$

 let $?O = (vector\ [0,\ 0])::real^{\wedge}2$
 let $?e2 = (vector\ [0,\ 1])::real^{\wedge}2$
 let $?y-unit-seg-path = linepath\ ?O\ ?e2$
 let $?y-unit-seg = path-image\ ?y-unit-seg-path$

 have $\forall t \in \{0..1\}. ?f\ (?y-unit-seg-path\ t) = p\ t$
proof *clarify*
 fix $t :: real$
 assume $t: t \in \{0..1\}$
 then obtain v where $v: ?y-unit-seg-path\ t = v$ by *auto*
 then have $v = (1 - t) *_{R}\ ?O + t *_{R}\ ?e2$ **unfolding** *linepath-def* by *auto*
 then have $v = t *_{R}\ ?e2$
 by (*smt* (*verit*, *best*) *t* *exhaust-2* *linepath-0* *scaleR-zero-left* *vec-eq-iff* *vector-2(1)* *vector-2(2)* *vector-scaleR-component*)
 then have $?f\ v = p\ t$
proof –
 assume $v = t *_{R}\ vector\ [0,\ 1]$
 then have $v = vector\ [t * 0,\ t * 1]$
 by (*smt* (*verit*, *del-insts*) *exhaust-2* *mult-cancel-left1* *real-scaleR-def* *scaleR-zero-right* *vec-eq-iff* *vector-2(1)* *vector-2(2)* *vector-scaleR-component*)
 then have $v: v = vector\ [0,\ t]$ by *auto*

 have $f1: ?f1\ v = vector\ [t * (b\ \$1 - a\ \$1),\ t * (b\ \$2 - a\ \$2)]$ (**is** $?f1\ v = ?f1-v$)
 by (*simp* *add: mat-vec-mult-2* v)

 have $?f2\ ?f1-v = vector\ [t * (b\ \$1 - a\ \$1),\ t * (b\ \$2 - a\ \$2)] + vector\ [a\ \$1,\ a\ \$2]$
 by (*smt* (*verit*) *exhaust-2* *vec-eq-iff* *vector-2(1)* *vector-2(2)*)
 also have $\dots = vector\ [t * (b\ \$1 - a\ \$1) + a\ \$1,\ t * (b\ \$2 - a\ \$2) + a\ \$2]$
 by (*smt* (*verit*, *del-insts*) *vector-add-component* *exhaust-2* *vec-eq-iff* *vec-*

```

tor-2(1) vector-2(2))
  also have ... = vector [t * b$1 + (1 - t) * a$1, t * b$2 + (1 - t) * a$2]
by argo
  also have ... = t *R b + (1 - t) *R a
    by (smt (verit, del-insts) exhaust-2 real-scaleR-def vec-eq-iff vector-2(1)
vector-2(2) vector-add-component vector-scaleR-component)
  finally have ?f2 ?f1-v = t *R b + (1 - t) *R a .
  thus ?thesis using p-def f1 unfolding linepath-def by simp
qed
thus ?f (?y-unit-seg-path t) = p t using v by simp
qed

then have ?f ' ?y-unit-seg = path-image p unfolding path-image-def by force
moreover have ?f differentiable-on ?y-unit-seg
proof-
  have linear ?f1 by auto
  then have ?f1 differentiable-on ?y-unit-seg
    using linear-imp-differentiable by (simp add: linear-imp-differentiable-on)
  moreover have ?f2 differentiable-on (?f1 ' ?y-unit-seg)
proof-
  have ?id differentiable-on ?f1 ' ?y-unit-seg
    using differentiable-const by simp
  moreover have ?f-a differentiable-on ?f1 ' ?y-unit-seg
    using differentiable-ident by simp
  ultimately show ?f2 differentiable-on ?f1 ' ?y-unit-seg
    using differentiable-compose by simp
qed
ultimately show ?thesis using differentiable-compose
  by (simp add: differentiable-chain-within differentiable-on-def)
qed
moreover have negligible ?y-unit-seg
  using linepath-is-negligible-vertical[of ?O ?e2] by simp
ultimately show ?thesis
  using negligible-differentiable-image-negligible by fastforce
qed

```

lemma linepath-is-negligible:

```

fixes a b :: real^2
defines p ≡ linepath a b
shows negligible (path-image p)
proof-
{ assume a$1 = b$1
  then have ?thesis using linepath-is-negligible-vertical p-def by blast
} moreover
{ assume a$1 < b$1
  then have ?thesis using linepath-is-negligible-non-vertical p-def by blast
} moreover
{ assume a: a$1 > b$1
  let ?p-rev = reversepath p

```

```

    have path-image p = path-image ?p-rev by simp
    moreover have ?p-rev = linepath b a using p-def by simp
    ultimately have ?thesis using a linepath-is-negligible-non-vertical[of b a] by
simp
  }
  ultimately show ?thesis by linarith
qed

```

```

lemma linepath-has-emeasure-0:
  emeasure lebesgue (path-image (linepath (a::(real^2)) (b::(real^2)))) = 0
  using linepath-is-negligible emeasure-notin-sets negligible-iff-emeasure0 by blast

```

```

lemma linepath-has-measure-0:
  measure lebesgue (path-image (linepath (a::(real^2)) (b::(real^2)))) = 0
  using linepath-has-emeasure-0 linepath-is-negligible negligible-imp-measure0 by
blast

```

```

end
theory Polygon-Convex-Lemmas
imports
  Polygon-Lemmas
  Linepath-Collinearity

```

```
begin
```

17 Misc. Convex Polygon Properties

```

lemma polygon-path-image-subset-convex:
  assumes length vts > 0
  shows path-image (make-polygonal-path vts)  $\subseteq$  convex hull (set vts) (is path-image
?p  $\subseteq$  ?S)
  using assms
proof(induct vts rule: make-polygonal-path.induct)
  case 1
  then show ?case by simp
next
  case (2 a)
  then show ?case by auto
next
  case (3 a b)
  show ?case (is path-image ?p  $\subseteq$  ?S)
  proof(rule subsetI)
    fix x
    assume x-in-path-image: x  $\in$  path-image ?p
    then have x  $\in$  path-image (linepath a b) by auto
    thus x  $\in$  ?S
  unfolding path-image-def linepath-def
  by (smt (verit, ccfv-SIG)  $\langle$ x  $\in$  path-image (linepath a b) $\rangle$  convex-alt con-
vex-convex-hull hull-subset in-mono in-segment(1) linepath-image-01 list.set-intros(1))

```

```

path-image-def set-subset-Cons)
qed
next
case (4 a b c tl)
let ?vts = a # b # c # tl
show ?case (is path-image ?p  $\subseteq$  ?S)
proof(rule subsetI)
  fix x
  assume x-in-path-image: x  $\in$  path-image ?p
  show x  $\in$  ?S
  proof cases
    assume x  $\in$  set ?vts
    thus ?thesis by (simp add: hull-inc)
  next
    assume x-notin: x  $\notin$  set ?vts
    obtain u where p-u: u  $\in$  {0..1}  $\wedge$  ?p u = x
      using x-in-path-image unfolding path-image-def by auto
    then have p-head-tail: ?p = (linepath a b) +++ make-polygonal-path (b # c
# tl)
      by auto
    have abc-in-S: set ?vts  $\subseteq$  convex hull (set ?vts) by (simp add: hull-subset)
    { assume u-assm: u  $\leq$  1/2
      then have ?p u = (1 - 2 * u) *R a + (2 * u) *R b
        using p-head-tail unfolding linepath-def joinpaths-def
        by presburger
      hence x  $\in$  ?S
        using abc-in-S convexD-alt[of ?S a b 2 * u] u-assm p-u by simp
    } moreover
    { assume u-assm: u > 1/2
      then have x = (make-polygonal-path (b # c # tl) (2 * u - 1)) (is x =
(?p' (2 * u - 1)))
        using p-head-tail p-u unfolding linepath-def joinpaths-def by auto
      moreover have 0 < (2 * u - 1) using u-assm by linarith
      ultimately have x  $\in$  path-image ?p'
        using p-u by (simp add: path-image-def)
      moreover have path-image ?p'  $\subseteq$  convex hull (set (b # c # tl)) using
4(1) by auto
      moreover have ...  $\subseteq$  convex hull (set (a # b # c # tl))
        by (meson hull-mono set-subset-Cons)
      ultimately have x  $\in$  ?S by auto
    }
  ultimately show ?thesis by linarith
qed
qed
qed

```

lemma *convex-contains-simple-closed-path-imp-contains-path-inside:*
 assumes *convex S*
 assumes *simple-path p \wedge closed-path p*

assumes *path-image* $p \subseteq S$
shows *path-inside* $p \subseteq S$
by (*metis* (*no-types*, *opaque-lifting*) *Compl-subset-Compl-iff Un-subset-iff* *assms(1)*
assms(3) *boolean-algebra-class.boolean-algebra.double-compl* *outside-subset-convex*
path-inside-def union-with-inside)

lemma *convex-polygon-is-convex-hull:*

assumes *polygon* p
assumes *convex* (*path-inside* $p \cup$ *path-image* p)
assumes $p =$ *make-polygonal-path* vts
shows *convex hull* (*set* vts) = *path-inside* $p \cup$ *path-image* p (**is** $?hull = ?poly$)
proof –
have $?hull \subseteq ?poly$
proof(*rule subsetI*)
fix x
assume $x \in ?hull$
moreover have $\forall H. (convex\ H \wedge (set\ vts) \subseteq H) \longrightarrow ?hull \subseteq H$ **by** (*simp*
add: hull-minimal)
moreover have *convex* ($?poly$) \wedge (*set* vts) \subseteq $?poly$
using *assms(2)* *assms(3)* *vertices-on-path-image* **by** *auto*
ultimately show $x \in ?poly$ **by** *auto*
qed
moreover have $?hull \supseteq ?poly$
proof(*rule subsetI*)
fix x
assume $x \in ?poly$
moreover have *path-image* $p \subseteq ?hull$
using *polygon-path-image-subset-convex[of vts]* *polygon-at-least-3-vertices*
assms
by *force*
moreover from *calculation* **have** *path-inside* $p \subseteq ?hull$
using *convex-contains-simple-closed-path-imp-contains-path-inside* *polygon-def*
assms(1)
by *auto*
ultimately show $x \in ?hull$ **by** *auto*
qed
ultimately show $?thesis$ **by** *auto*
qed

lemma *convex-polygon-inside-is-convex-hull-interior:*

assumes *polygon* p
assumes *convex* (*path-inside* p)
assumes $p =$ *make-polygonal-path* vts
shows *interior* (*convex hull* (*set* vts)) = *path-inside* p
by (*metis* (*no-types*, *lifting*) *assms closure-Un-frontier* *convex-closure* *convex-interior-closure*
convex-polygon-is-convex-hull *inside-outside-def* *inside-outside-polygon* *interior-eq*)

lemma *convex-polygon-inside-is-convex-hull-interior2:*

assumes *polygon* p

assumes *convex* (*path-inside* $p \cup \text{path-image } p$)
assumes $p = \text{make-polygonal-path } vts$
shows *interior* (*convex hull* (*set* vts)) = *path-inside* p
using *assms* *closure-Un-frontier* *convex-closure* *convex-interior-closure* *convex-polygon-is-convex-hull*
inside-outside-def *inside-outside-polygon* *interior-eq*
by (*smt* (*verit*, *best*) *List.finite-set* *compact-eq-bounded-closed* *finite-imp-compact-convex-hull*
frontier-complement *inside-frontier-eq-interior* *outside-inside* *path-inside-def* *path-outside-def*
sup-commute)

lemma *polygon-convex-iff*:
assumes *polygon* p
shows *convex* (*path-inside* p) \longleftrightarrow *convex* (*path-inside* $p \cup \text{path-image } p$)
using *convex-polygon-inside-is-convex-hull-interior*
using *convex-polygon-inside-is-convex-hull-interior2*
by (*metis* *Jordan-inside-outside-real2* *closed-path-def* *assms* *closure-Un-frontier*
convex-closure *convex-interior* *convex-polygon-is-convex-hull* *path-inside-def* *poly-*
gon-def *polygon-to-polygonal-path*)

lemma *convex-polygon-frontier-is-path-image*:
assumes *polygon-of* p vts
assumes *convex* (*path-inside* p)
shows *frontier* (*convex hull* (*set* vts)) = *path-image* p
using *assms*
unfolding *frontier-def* *polygon-of-def*
by (*metis* (*no-types*, *lifting*) *Jordan-inside-outside-real2* *closed-path-def* *convex-closure-interior*
convex-convex-hull *convex-polygon-inside-is-convex-hull-interior* *frontier-def* *inter-*
rior-interior *path-inside-def* *polygon-def*)

lemma *convex-polygon-frontier-is-path-image2*:
assumes *polygon* p
assumes *convex* (*path-inside* p)
shows *frontier* (*path-image* $p \cup \text{path-inside } p$) = *path-image* p
using *assms*
by (*simp* *add*: *Jordan-inside-outside-real2* *closed-path-def* *path-inside-def* *poly-*
gon-def *union-with-inside*)

lemma *convex-polygon-frontier-is-path-image3*:
assumes *polygon* p
assumes *convex* (*path-image* $p \cup \text{path-inside } p$)
shows *frontier* (*path-image* $p \cup \text{path-inside } p$) = *path-image* p
using *assms* *polygon-convex-iff*
by (*simp* *add*: *convex-polygon-frontier-is-path-image2* *sup-commute*)

lemma *polygon-frontier-is-path-image*:
assumes *polygon* p
shows *frontier* (*path-inside* p) = *path-image* p
using *inside-outside-polygon* **unfolding** *inside-outside-def*
using *assms* **by** *presburger*

lemma *convex-path-inside-means-convex-polygon*:
assumes *polygon* p
assumes *frontier* $(\text{convex hull } (\text{set } vts)) = \text{path-image } p$
shows *convex* $(\text{path-inside } p)$
by $(\text{metis List.finite-set assms}(2) \text{convex-convex-hull convex-interior finite-imp-bounded-convex-hull inside-frontier-eq-interior path-inside-def})$

lemma *convex-hull-of-polygon-is-convex-hull-of-vts*:
assumes *polygon-of* p vts
shows $\text{convex hull } (\text{path-image } p \cup \text{path-inside } p) = \text{convex hull } (\text{set } vts)$
proof –
have *len-vts*: $\text{length } vts > 0$
by $(\text{metis assms card.empty empty-set length-greater-0-conv not-numeral-le-zero polygon-at-least-3-vertices polygon-of-def})$
have $\text{path-image } p \cup \text{path-inside } p \subseteq \text{convex hull } (\text{set } vts)$
using *polygon-path-image-subset-convex* $[OF \text{ len-vts}]$
using *assms convex-contains-simple-closed-path-imp-contains-path-inside polygon-def polygon-of-def* **by** *auto*
then have *subset1*: $\text{convex hull } (\text{path-image } p \cup \text{path-inside } p) \subseteq \text{convex hull } (\text{set } vts)$
by $(\text{simp add: convex-hull-subset})$
have $\text{set } vts \subseteq \text{path-image } p \cup \text{path-inside } p$ **using** *assms vertices-on-path-image*

by $(\text{simp add: polygon-of-def sup.coboundedI1})$
then have *subset2*: $\text{convex hull } (\text{set } vts) \subseteq \text{convex hull } (\text{path-image } p \cup \text{path-inside } p)$
by $(\text{simp add: hull-mono})$
show *?thesis* **using** *subset1 subset2*
by *auto*
qed

lemma *convex-hull-frontier-polygon*:
assumes *polygon-of* p vts
assumes $\neg \text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\neg \text{convex } (\text{path-inside } p)$
by $(\text{metis assms}(1) \text{assms}(2) \text{convex-polygon-frontier-is-path-image polygon-of-def vertices-on-path-image})$

lemma *frontier-int-subset*:
assumes $A \subseteq B$
shows $(\text{frontier } B) \cap A \subseteq \text{frontier } A$
by $(\text{metis assms closure-Un-frontier frontier-Int inf.absorb-iff2 inf-sup-aci}(1) \text{subset-Un-eq sup-inf-distrib2})$

lemma *in-frontier-in-subset*:
assumes $A \subseteq B$
assumes $x \in \text{frontier } B$
assumes $x \in A$
shows $x \in \text{frontier } A$

by (metis assms frontier-int-subset IntI in-mono)

lemma *in-frontier-in-subset-convex-hull*:

assumes $A \subseteq B$
assumes $x \in \text{frontier (convex hull } B)$
assumes $x \in \text{convex hull } A$
shows $x \in \text{frontier (convex hull } A)$
by (metis in-frontier-in-subset assms hull-mono)

lemma *convex-hull-two-extreme-points*:

fixes $S :: 'a::\text{euclidean-space set}$
assumes *finite* S
assumes $\text{convex hull } S \neq \{\}$
assumes $\forall x. \text{convex hull } S \neq \{x\}$
shows $\text{card } \{x. x \text{ extreme-point-of (convex hull } S)\} \geq 2$ (is $\text{card } ?ep \geq 2$)

proof –

have *compact (convex hull S)* by (simp add: assms(1) *finite-imp-compact-convex-hull*)
then have $\text{convex hull } S = \text{convex hull } ?ep$

using *Krein-Milman-Minkowski[OF - convex-convex-hull]* by blast
moreover then obtain x where $x \in ?ep$ using assms(2) by fastforce
moreover have $?ep \neq \{x\}$ using assms(3) *calculation(1)* by force
ultimately obtain y where $x \in ?ep \wedge y \in ?ep \wedge x \neq y$ by blast

moreover have *finite ?ep* using assms(1) *extreme-points-of-convex-hull finite-subset*
by blast

ultimately show *?thesis*

by (metis (no-types, lifting) *One-nat-def Orderings.order-eq-iff Suc-1 Suc-leI card-1-singletonE card-gt-0-iff empty-iff insert-Diff not-less-eq-eq singleton-insert-inj-eq*)
qed

lemma *convex-hull-two-pts-on-frontier*:

fixes $S :: 'a::\text{euclidean-space set}$
assumes $\text{card } S \geq 2$
shows $\text{card } (S \cap \text{frontier (convex hull } S)) \geq 2$

proof –

have $S \subseteq \text{convex hull } S$ by (simp add: *hull-subset*)
then have $\text{convex hull } S \neq \{\} \wedge \text{card (convex hull } S) \neq 1$

by (metis *Suc-1 add-leD2 assms card.empty card-1-singletonE convex-hull-eq-empty not-one-le-zero numeral-le-one-iff plus-1-eq-Suc semiring-norm(69) subset-singletonD*)
moreover have *finite S* using assms by (metis *Suc-1 Suc-leD card-eq-0-iff not-one-le-zero*)

ultimately have $\text{card } \{x. x \text{ extreme-point-of (convex hull } S)\} \geq 2$

using *convex-hull-two-extreme-points* by fastforce

moreover have $\{x. x \text{ extreme-point-of (convex hull } S)\} \subseteq S \cap \text{frontier (convex hull } S)$

proof –

have $\{x. x \text{ extreme-point-of (convex hull } S)\} \subseteq S$ by (simp add: *extreme-points-of-convex-hull*)

moreover have $\{x. x \text{ extreme-point-of (convex hull } S)\} \cap \text{interior (convex hull } S) = \{\}$

using *extreme-point-not-in-interior* by blast

moreover have $\{x. x \text{ extreme-point-of } (\text{convex hull } S)\} \subseteq \text{convex hull } S$
using $\langle S \subseteq \text{convex hull } S \rangle \text{ calculation}(1)$ **by** *blast*
moreover have $\text{convex hull } S = \text{interior } (\text{convex hull } S) \cup \text{frontier } (\text{convex hull } S)$
by (*metis (no-types, lifting) Diff-empty Suc-1 assms card.infinite closure-Un-frontier closure-convex-hull convex-closure-interior convex-convex-hull empty-subsetI finite-imp-compact frontier-def interior-interior not-less-eq-eq sup-absorb2 zero-less-one-class.zero-le-one*)
ultimately show *?thesis* **by** *blast*
qed
ultimately show *?thesis*
by (*smt (verit, del-insts) assms extreme-points-of-convex-hull card-gt-0-iff finite-Int linorder-not-less not-numeral-le-zero order-less-le order-less-le-trans psubset-card-mono*)
qed

18 Vertices on Convex Frontier Implies Polygon is Convex

lemma *convex-cut-aux*:

assumes $\forall v \in S. z \cdot v \leq 0$
shows $\text{convex hull } S \subseteq \{x. z \cdot x \leq 0\}$
by (*simp add: assms convex-halfspace-le hull-minimal subsetI*)

lemma *convex-cut-aux'*:

assumes $\forall v \in S. z \cdot v \geq 0$
shows $\text{convex hull } S \subseteq \{x. z \cdot x \geq 0\}$
using *convex-cut-aux[of S -z] assms* **by** *auto*

lemma *convex-cut*:

assumes $z \neq 0$
assumes $\{x. z \cdot x = 0\} \cap \text{interior } (\text{convex hull } S) \neq \{\}$
obtains $v1\ v2$ **where** $v1 \neq v2 \wedge \{v1, v2\} \subseteq S \wedge v1 \in \{x. z \cdot x < 0\} \wedge v2 \in \{x. z \cdot x > 0\}$
proof –
let $?P1 = \{x. z \cdot x \leq 0\}$
let $?P2 = \{x. z \cdot x \geq 0\}$
have $\text{frontier } ?P1 = \{x. z \cdot x = 0\}$
by (*simp add: assms(1) frontier-halfspace-le*)
moreover have $\text{frontier } ?P2 = \{x. z \cdot x = 0\}$
by (*simp add: assms(1) frontier-halfspace-ge*)
ultimately have $\neg \text{convex hull } S \subseteq ?P1 \wedge \neg \text{convex hull } S \subseteq ?P2$
by (*smt (verit, ccfv-SIG) DiffE IntE assms(2) disjoint-iff frontier-def inf.absorb-iff2 interior-Int*)
moreover have $(\forall v \in S. z \cdot v \leq 0) \implies \text{convex hull } S \subseteq ?P1$ **using** *convex-cut-aux* **by** *blast*
moreover have $(\forall v \in S. z \cdot v \geq 0) \implies \text{convex hull } S \subseteq ?P2$ **using** *convex-cut-aux'* **by** *blast*
ultimately obtain $v1\ v2$ **where** $\{v1, v2\} \subseteq S \wedge z \cdot v1 < 0 \wedge z \cdot v2 > 0$

using *linorder-not-le* **by** *auto*
thus *?thesis using that* **by** *fastforce*
qed

lemma *affine-2-int-convex*:

fixes $S :: 'a::\text{euclidean-space set}$
assumes $\{a, b\} \subseteq S$
assumes $\{a, b\} \subseteq \text{frontier } (\text{convex hull } S)$
assumes $\text{affine hull } \{a, b\} \cap \text{interior } (\text{convex hull } S) \neq \{\}$
shows $\text{affine hull } \{a, b\} \cap \text{convex hull } S = \text{convex hull } \{a, b\}$
proof –
let $?H = \text{convex hull } S$
let $?L = \text{affine hull } \{a, b\} \cap ?H$
have $1: ?L \supseteq \text{convex hull } \{a, b\}$
by (*meson Int-greatest assms(1) convex-hull-subset-affine-hull hull-mono*)
moreover **have** $?L \subseteq \text{convex hull } \{a, b\}$
proof(*rule subsetI*)
fix x
assume $*$: $x \in ?L$
then obtain $u v$ **where** $uv: x = u *_R a + v *_R b \wedge u + v = 1$ **using**
affine-hull-2 **by** *blast*

have $\text{rel-interior } ?L \subseteq \text{rel-interior } ?H$
using *subset-rel-interior-convex[of ?L ?H]*
by (*metis assms(3) convex-affine-hull convex-convex-hull convex-rel-interior-inter-two inf-bot-right inf-le2 rel-interior-affine-hull rel-interior-nonempty-interior*)
moreover **have** $ab\text{-frontier}: a \in \text{frontier } ?H \wedge b \in \text{frontier } ?H$ **using** *assms*
by *blast*
ultimately **have** $ab\text{-rel-frontier}: a \in \text{rel-frontier } ?L \wedge b \in \text{rel-frontier } ?L$
by (*metis IntI affine-affine-hull assms(3) convex-affine-rel-frontier-Int convex-convex-hull hull-subset inf-commute insert-subset*)

{ assume $**$: $u < 0$
then **have** $b \in \text{open-segment } a x$
proof –
from uv **have** $b = (1/v) *_R x - (u/v) *_R a$
by (*smt (verit, ccfv-threshold) ** divide-inverse-commute inverse-eq-divide real-vector-affinity-eq vector-space-assms(3) Groups.add-ac(2)*)
moreover **from** uv **have** $1/v - u/v = 1$
by (*metis ** add.commute add-cancel-right-left diff-divide-distrib divide-self-if eq-diff-eq! not-one-less-zero*)
ultimately **have** $b = (1 - 1/v) *_R a + (1/v) *_R x$ **by** (*simp add: diff-eq-eq*)
moreover **from** $uv **$ **have** $0 < 1/v \wedge 1/v < 1$ **by** *simp*
ultimately **show** *?thesis*
by (*metis 1 ab-rel-frontier affine-hull-sing convex-hull-singleton empty-iff equalityI in-segment(2) inf-le1 insert-absorb rel-frontier-sing scaleR-collapse singletonI*)
qed
then **have** $b \in \text{rel-interior } (\text{convex hull } \{a, x\})$

by (*metis empty-iff open-segment-idem rel-interior-closed-segment segment-convex-hull*)
moreover have $x \in ?H$ **using** * **by** *blast*
ultimately have $b \in \text{interior } ?H$
by (*smt (verit, ccfv-threshold) * IntD2 Int-empty-right 1 affine-affine-hull affine-hull-affine-Int-nonempty-interior affine-hull-convex-hull assms(3) convex-Int convex-affine-hull convex-convex-hull convex-rel-interior-inter-two hull-hull hull-redundant-eq insert-commute insert-subsetI rel-interior-affine-hull rel-interior-mono rel-interior-nonempty-interior rel-interior-subset subset-hull subset-iff*)
then have *False* **by** (*metis DiffD2 ab-frontier frontier-def*)
} **moreover**
{ **assume** **: $v < 0$
then have $a \in \text{open-segment } b \ x$
proof-
from uv **have** $a = (1/u) *_R x - (v/u) *_R b$
by (*smt (verit, ccfv-threshold) ** divide-inverse-commute inverse-eq-divide real-vector-affinity-eq vector-space-assms(3) Groups.add-ac(2)*)
moreover from uv **have** $1/u - v/u = 1$
by (*metis ** add-cancel-right-left diff-divide-distrib divide-self-if eq-diff-eq' not-one-less-zero*)
ultimately have $a = (1 - 1/u) *_R b + (1/u) *_R x$ **by** (*simp add: diff-eq-eq*)
moreover from uv **** have** $0 < 1/u \wedge 1/u < 1$ **by** *simp*
ultimately show *?thesis*
by (*metis 1 ab-rel-frontier affine-hull-sing convex-hull-singleton empty-iff equalityI in-segment(2) inf-le1 insert-absorb rel-frontier-sing scaleR-collapse singletonI*)
qed
then have $a \in \text{rel-interior } (\text{convex hull } \{b, x\})$
by (*metis empty-iff open-segment-idem rel-interior-closed-segment segment-convex-hull*)
moreover have $x \in ?H$ **using** * **by** *blast*
ultimately have $a \in \text{interior } ?H$
by (*smt (verit, ccfv-threshold) * IntD2 Int-empty-right 1 affine-affine-hull affine-hull-affine-Int-nonempty-interior affine-hull-convex-hull assms(3) convex-Int convex-affine-hull convex-convex-hull convex-rel-interior-inter-two hull-hull hull-redundant-eq insert-commute insert-subsetI rel-interior-affine-hull rel-interior-mono rel-interior-nonempty-interior rel-interior-subset subset-hull subset-iff*)
then have *False* **by** (*metis DiffD2 ab-frontier frontier-def*)
}
ultimately have $0 \leq u \wedge u \leq 1 \wedge 0 \leq v \wedge v \leq 1$ **using** uv **by** *argo*
thus $x \in \text{convex hull } \{a, b\}$ **by** (*simp add: convexD hull-inc uv*)
qed
ultimately show *?thesis* **by** *blast*
qed

lemma *halfplane-frontier-affine-hull:*

fixes $b \ v :: \text{real}^2$
assumes $b \neq 0$
assumes $v \neq 0$

```

assumes  $b \in \{x. v \cdot x = 0\}$ 
shows  $\{x. v \cdot x = 0\} = \text{affine hull } \{0, b\}$ 
proof –
  let  $?F = \{x. v \cdot x = 0\}$ 
  let  $?A = \text{affine hull } \{0, b\}$ 
  have  $?F \subseteq ?A$ 
  proof(rule subsetI)
    fix  $y$ 
    assume  $*$ :  $y \in ?F$ 
    have  $y \in ?A$  if  $y = 0$  by (simp add: assms(2) hull-inc that)
    moreover have  $y \in ?A$  if  $b \neq 0$ 
    proof–
      have  $v \cdot y = 0$  using  $*$  by fast
      moreover have  $v \cdot b = 0$  using assms by force
      moreover have  $v \cdot y = v_1 * y_1 + v_2 * y_2$  by (simp add: inner-vec-def sum-2 real-2-inner)
      moreover have  $v \cdot b = v_1 * b_1 + v_2 * b_2$  by (simp add: inner-vec-def sum-2 real-2-inner)
      ultimately have  $0: v_1 * y_1 + v_2 * y_2 = 0 \wedge 0 = v_1 * b_1 + v_2 * b_2$  by auto
      moreover obtain  $c$  where  $c: y_1 = c * b_1$  using  $\langle b_1 \neq 0 \rangle$ 
        by (metis hyperplane-eq-Ex inner-real-def mult.commute)
      ultimately have  $v_1 * y_1 + v_2 * y_2 = 0 \wedge 0 = c * v_1 * b_1 + c * v_2 * b_2$  by algebra
      then have  $v_1 * y_1 + v_2 * y_2 = v_1 * y_1 + c * v_2 * b_2$  using  $c$  by algebra
      then have  $v_2 * y_2 = c * v_2 * b_2$  by argo
      then have  $y_2 = c * b_2$ 
        by (smt (verit, ccfv-threshold) 0 exhaust-2 mult.commute mult.left-commute mult-cancel-left that assms vec-eq-iff zero-index)
      then have  $y = c *_{\mathbb{R}} b$  using  $c$ 
        by (smt (verit) exhaust-2 real-scaleR-def vec-eq-iff vector-scaleR-component)
      then have  $y \in \text{span } \{0, b\}$  by (meson insert-subset span-mul span-superset)
      thus  $y \in ?A$ 
        by (simp add: affine-hull-span-0 assms(2) hull-inc)
    qed
  moreover have  $y \in ?A$  if  $b_2 \neq 0$ 
  proof–
    have  $v \cdot y = 0$  using  $*$  by fast
    moreover have  $v \cdot b = 0$  using assms by force
    moreover have  $v \cdot y = v_1 * y_1 + v_2 * y_2$  by (simp add: inner-vec-def sum-2 real-2-inner)
    moreover have  $v \cdot b = v_1 * b_1 + v_2 * b_2$  by (simp add: inner-vec-def sum-2 real-2-inner)
    ultimately have  $0: v_1 * y_1 + v_2 * y_2 = 0 \wedge 0 = v_1 * b_1 + v_2 * b_2$  by auto
    moreover obtain  $c$  where  $c: y_2 = c * b_2$  using  $\langle b_2 \neq 0 \rangle$ 
      by (metis hyperplane-eq-Ex inner-real-def mult.commute)
    ultimately have  $v_1 * y_1 + v_2 * y_2 = 0 \wedge 0 = c * v_1 * b_1 + c * v_2 * b_2$ 

```

$v\$2 * b\2 by algebra
then have $v\$1 * y\$1 + v\$2 * y\$2 = 0 \wedge 0 = c * v\$1 * b\$1 + v\$2 * y\2
using c by algebra
then have $v\$1 * y\$1 = c * v\$1 * b\1 by argo
then have $y\$1 = c * b\1
by (smt (verit, ccfv-threshold) 0 exhaust-2 mult.commute mult.left-commute mult-cancel-left that assms vec-eq-iff zero-index)
then have $y = c *_R b$ using c
by (smt (verit) exhaust-2 real-scaleR-def vec-eq-iff vector-scaleR-component)
then have $y \in \text{span } \{0, b\}$ by (meson insert-subset span-mul span-superset)
thus $y \in ?A$
by (simp add: affine-hull-span-0 assms(2) hull-inc)
qed
ultimately show $y \in ?A$
by (metis (mono-tags, opaque-lifting) assms(1) exhaust-2 vec-eq-iff zero-index)
qed
moreover have $?A \subseteq ?F$
proof(rule subsetI)
fix x
assume $x \in ?A$
then obtain $\alpha \beta$ **where** $x = \alpha *_R 0 + \beta *_R b \wedge \alpha + \beta = 1$ using affine-hull-2
by blast
then have $v \cdot x = \alpha * (v \cdot 0) + \beta * (v \cdot b)$ by (simp add: assms(1))
then have $v \cdot x = 0$ using assms(3) by auto
thus $x \in ?F$ by fast
qed
ultimately show $?thesis$ by blast
qed

lemma vts-on-convex-frontier-aux:

assumes polygon-of p vts
assumes $vts!0 = 0$
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{path-image } (\text{linepath } (vts!0) (vts!1)) \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$

proof–

let $?H = \text{convex hull } (\text{set } vts)$
let $?a = vts!0$
let $?b = vts!1$
let $?l = \text{linepath } ?a ?b$
let $?L = \text{path-image } ?l$
let $?A = \text{affine hull } \{?a, ?b\}$
let $?x = ?b - ?a$

obtain v **where** $v \cdot ?x = 0 \wedge v \neq 0$

proof–

let $?v = (\text{vector } [?x\$2, -?x\$1])::(\text{real}^2)$
have $?a \neq ?b$
by (smt (verit, best) Cons-nth-drop-Suc One-nat-def Suc-le-eq arc-distinct-ends assms(1) assms(2) card.empty drop0 empty-set length-greater-0-conv list.sel(1))

```

list.sel(3) make-polygonal-path.elims make-polygonal-path.simps(1) make-polygonal-path.simps(2)
nth-drop pathfinish-linepath pathstart-linepath plus-1-eq-Suc polygon-at-least-3-vertices
polygon-def polygon-of-def polygon-pathstart rel-simps(28) simple-path-joinE)
  then have ?x ≠ 0 by simp
  then have ?v · ?x = 0 ∧ ?v ≠ 0
  proof-
    have ?v · ?x = (?x$2 * ?x$1) + (-?x$1 * ?x$2)
      by (simp add: inner-vec-def sum-2 real-2-inner)
    then have ?v · ?x = 0 by argo
    moreover have ?v ≠ 0
      by (smt (verit, best) ⟨?x ≠ 0⟩ exhaust-2 vec-eq-iff vector-2(1) vector-2(2)
zero-index)
    ultimately show ?thesis by blast
  qed
  thus ?thesis using that by blast
qed

let ?P1 = {x. v · x ≤ 0}
let ?P2 = {x. v · x ≥ 0}
let ?P1-int = {x. v · x < 0}
let ?P2-int = {x. v · x > 0}
let ?F = {x. v · x = 0}

have ?b ≠ 0
  by (smt (verit) Cons-nth-drop-Suc One-nat-def Suc-le-eq Suc-le-length-iff arc-distinct-ends
assms(1) assms(2) card.empty drop0 drop-eq-Nil empty-set le-numeral-extra(4)
length-greater-0-conv list.inject make-polygonal-path.elims make-polygonal-path.simps(2)
nat-less-le pathfinish-linepath pathstart-linepath polygon-at-least-3-vertices polygon-def
polygon-of-def polygon-pathstart rel-simps(28) simple-path-joinE)
  moreover have ?b ∈ ?F using assms(2) v by auto
  ultimately have F: ?F = ?A
    using halfplane-frontier-affine-hull[of ?b v] v assms(2) by presburger
  moreover have ?L ⊆ ?A by (simp add: convex-hull-subset-affine-hull segment-convex-hull)
  ultimately have L-subset-F: ?L ⊆ ?F by blast
  have L-subset-H: ?L ⊆ ?H
    by (metis (no-types, lifting) add-gr-0 assms(1) card.empty convex-contains-segment
convex-convex-hull diff-less empty-set hull-subset leD length-greater-0-conv less-numeral-extra(1)
nth-mem numeral-3-eq-3 path-image-linepath plus-1-eq-Suc polygon-at-least-3-vertices
polygon-of-def rotate-polygon-vertices-same-set rotated-polygon-vertices-helper(2) sub-
set-code(1))

have frontier-P1: frontier ?P1 = ?F by (simp add: v frontier-halfspace-le)
have frontier-P2: frontier ?P2 = ?F by (simp add: v frontier-halfspace-ge)
have interior-P1: interior ?P1 = ?P1-int by (simp add: v)
have interior-P2: interior ?P2 = ?P2-int by (simp add: v)
have convex-P1: convex ?P1 by (simp add: convex-halfspace-le)
have convex-P2: convex ?P2 by (simp add: convex-halfspace-ge)
have P1-int-P2: ?P1 ∩ ?P2 = ?F by (simp add: halfspace-Int-eq(1))

```

```

let ?H1 = ?H ∩ ?P1
let ?H2 = ?H ∩ ?P2

have ¬ collinear (set vts) using polygon-vts-not-collinear assms(1) by simp
then have nonempty-interior-H: interior ?H ≠ {}
  by (smt (verit, ccfv-SIG) Jordan-inside-outside-real2 closed-path-def Un-Int-eq(4)
  assms(1) convex-hull-of-polygon-is-convex-hull-of-vts disjoint-iff hull-subset inf.orderE
  interior-Int interior-eq interior-subset path-inside-def polygon-def polygon-of-def)

have convex-H1: convex ?H1 by (simp add: convex-Int convex-P1)
have convex-H2: convex ?H2 by (simp add: convex-Int convex-P2)

have ?H ⊆ ?P1 ∨ ?H ⊆ ?P2
proof(rule ccontr)
  assume *: ¬ (?H ⊆ ?P1 ∨ ?H ⊆ ?P2)
  moreover have interior ?H ⊆ ?P1 ⇒ ?H ⊆ ?P1
    by (metis (no-types, lifting) Int-Un-eq(3) Krein-Milman-frontier List.finite-set
    P1-int-P2 closure-Un-frontier closure-convex-hull closure-mono compact-frontier con-
    vex-closure-interior convex-convex-hull finite-imp-compact-convex-hull frontier-P1
    nonempty-interior-H)
  moreover have interior ?H ⊆ ?P2 ⇒ ?H ⊆ ?P2
    by (metis (no-types, lifting) Int-Un-eq(3) Krein-Milman-frontier List.finite-set
    P1-int-P2 calculation(1) calculation(2) closure-Un-frontier closure-convex-hull clo-
    sure-mono compact-frontier convex-closure-interior convex-convex-hull emptyE fi-
    nite-imp-compact-convex-hull frontier-P2 inf-commute subsetI)
  ultimately have interior ?H ∩ ?P1 ≠ {} ∧ interior ?H ∩ ¬?P1 ≠ {} by
  force
  moreover have path-connected (interior ?H) by (simp add: convex-imp-path-connected)
  ultimately have F-int-interior-H: ?F ∩ interior ?H ≠ {}
  by (metis (no-types, lifting) path-connected-frontier ComplD disjoint-eq-subset-Compl
  frontier-P1 subset-eq)
  then obtain v1 v2 where v1v2: v1 ≠ v2 ∧ {v1, v2} ⊆ set vts
    ∧ v1 ∈ interior ?P1 ∧ v2 ∈ interior ?P2
  using convex-cut frontier-P1 interior-P1 interior-P2 v by metis
  then obtain i j where ij: vts!i = v1 ∧ vts!j = v2
    ∧ 2 ≤ i ∧ 2 ≤ j ∧ i ≠ j ∧ i < length vts - 1 ∧ j < length vts - 1
  proof-
    obtain i j where vts!i = v1 ∧ vts!j = v2 ∧ i ≠ j ∧ i < length vts ∧ j <
    length vts
    by (metis in-set-conv-nth insert-subset v1v2)
    moreover have 2 ≤ i
    proof-
      { assume i = 0 ∨ i = 1
        then have vts!i = ?a ∨ vts!i = ?b by blast
        then have vts!i ∈ ?F by (simp add: F hull-inc)
        then have False using calculation(1) interior-P1 v1v2 by auto
        }
    thus ?thesis by presburger
  qed

```

```

moreover have  $2 \leq j$ 
proof-
  { assume  $j = 0 \vee j = 1$ 
    then have  $vts!j = ?a \vee vts!j = ?b$  by blast
    then have  $vts!j \in ?F$  by (simp add: F hull-inc)
    then have False using calculation(1) interior-P2 v1v2 by auto
  }
thus ?thesis by presburger
qed
moreover have False if  $i = \text{length } vts - 1$ 
by (metis (no-types, lifting) F assms(1) calculation(1) frontier-P1 frontier-def
have-wraparound-vertex hull-subset insertCI insert-Diff last-conv-nth last-snoc less-nat-zero-code
list.size(3) polygon-of-def subset-Diff-insert that v1v2)
moreover have False if  $j = \text{length } vts - 1$ 
by (metis (no-types, lifting) F assms(1) calculation(1) frontier-P2 frontier-def
have-wraparound-vertex hull-subset insertCI insert-Diff last-conv-nth last-snoc less-nat-zero-code
list.size(3) polygon-of-def subset-Diff-insert that v1v2)
ultimately show ?thesis using that by fastforce
qed

let  $?i' = \min i j$ 
let  $?j' = \max i j$ 
let  $?vts' = \text{take } (?j' - ?i' + 1) (\text{drop } ?i' vts)$ 
let  $?p' = \text{make-polygonal-path } ?vts'$ 
have  $vts'\text{-sublist: sublist } ?vts' vts$  using sublist-order.order.trans by blast
then have  $vts'\text{-sublist-tl: sublist } ?vts' (\text{tl } vts)$ 
by (metis Suc-1 Suc-eq-plus1 drop-Suc ij max-def min-def nat-minus-add-max
not-less-eq-eq sublist-drop sublist-order.dual-order.trans sublist-take)

have  $p'\text{-start-finish: } \{\text{pathstart } ?p', \text{pathfinish } ?p'\} = \{v1, v2\}$ 
proof-
  have  $?vts'!0 = vts'! ?i'$  using ij by force
  moreover have  $?vts'!(?j' - ?i') = vts'! ?j'$ 
  using diff-is-0-eq diff-zero ij less-numeral-extra(1) max.cobounded1 min-absorb2
min-def nth-drop nth-take order-less-imp-le
  by fastforce
  moreover have  $(vts'! ?i' = v1 \wedge vts'! ?j' = v2) \vee (vts'! ?i' = v2 \wedge vts'! ?j' = v1)$ 
  using ij by linarith
  moreover have  $\text{pathstart } ?p' = ?vts'!0 \wedge \text{pathfinish } ?p' = ?vts'!(?j' - ?i')$ 
  using ij min-diff polygon-pathfinish polygon-pathstart
  by (smt (verit, ccfv-SIG) add-diff-cancel-right' add-diff-inverse-nat length-drop
length-take less-diff-conv max commute max-min-same(1) min.absorb4 nat-minus-add-max
not-add-less2 plus-1-eq-Suc plus-nat.simps(2) take-eq-Nil zero-less-one)
  ultimately show ?thesis by auto
qed
then have  $\text{path-image } ?p' \cap \text{interior } ?P2 \neq \{\} \wedge \text{path-image } ?p' \cap \text{interior } ?P1 \neq \{\}$ 
by (metis v1v2 IntI doubleton-eq-iff empty-iff pathfinish-in-path-image path-start-in-path-image)

```

then have $\text{path-image } ?p' \cap -?P1 \neq \{\} \wedge \text{path-image } ?p' \cap ?P1 \neq \{\}$
using *interior-P2*
by (*smt (verit, best) disjoint-iff-not-equal in-mono inf-shunt interior-P1 mem-Collect-eq*)
moreover have $\text{path-connected } (\text{path-image } ?p')$
using *make-polygonal-path-gives-path path-connected-path-image* **by** *blast*
ultimately obtain z **where** $z: z \in \text{path-image } ?p' \cap ?F$
by (*smt (verit, del-insts) path-connected-frontier DiffE Diff-triv all-not-in-conv frontier-P1*)
moreover have $\text{path-image } ?p' \subseteq ?H$
proof-
have $\text{path-image } p \subseteq ?H$
by (*metis assms(1) insert-subset length-pos-if-in-set polygon-of-def polygon-path-image-subset-convex v1v2*)
moreover have $\text{path-image } ?p' \subseteq \text{path-image } p$
by (*metis (no-types, lifting) vts'-sublist sublist-path-image-subset One-nat-def Suc-leI p'-start-finish assms(1) doubleton-eq-iff length-greater-0-conv make-polygonal-path.simps(1) pathfinish-linepath pathstart-linepath polygon-of-def v1v2*)
ultimately show *?thesis* **by** *blast*
qed
ultimately have $z \in \text{path-image } ?p' \cap (?H \cap ?F)$ **by** *blast*
moreover have $?H \cap ?F = ?L$
using *affine-2-int-convex[of ?a ?b set vts]*
by (*smt (verit, best) assms(3) F F-int-interior-H inf-commute segment-convex-hull path-image-linepath Suc-1 add-leD2 assms(1) empty-subsetI insert-subset length-greater-0-conv lessI nat-neq-iff nth-mem numeral-Bit0 order.strict-iff-not plus-1-eq-Suc polygon-of-def polygon-vertices-length-at-least-4 take-all-iff take-eq-Nil IntE inf.orderE*)
ultimately have $z \in ?L \cap \text{path-image } ?p'$ **by** *blast*
moreover have $?L \cap \text{path-image } ?p' \subseteq \{?a, ?b\}$
proof-
let $?p\text{-tl} = \text{make-polygonal-path } (tl \ vts)$
have $p = \text{make-polygonal-path } vts \wedge \text{loop-free } p$
using *assms unfolding polygon-of-def polygon-def simple-path-def* **by** *blast*
moreover have $[?a, ?b] = \text{take } 2 \ vts$
by (*metis Cons-nth-drop-Suc One-nat-def Suc-1 append-Cons append-Nil calculation constant-linepath-is-not-loop-free drop0 drop-eq-Nil insert-subset length-pos-if-in-set linorder-not-le make-polygonal-path.simps(2) take0 take-Suc-conv-app-nth v1v2*)
moreover have $tl \ vts = \text{drop } (2 - 1) \ vts$ **by** (*simp add: drop-Suc*)
moreover have $?l = \text{make-polygonal-path } [?a, ?b]$ **using** *make-polygonal-path.simps*
by *simp*
moreover have $\text{length } vts > 2$ **using** *ij* **by** *linarith*
moreover have $\text{pathstart } ?l = ?a \wedge \text{pathstart } ?p\text{-tl} = ?b$
using *calculation(3) calculation(5) polygon-pathstart* **by** *auto*
ultimately have $?L \cap \text{path-image } ?p\text{-tl} \subseteq \{?a, ?b\}$
using *loop-free-split-int[of p vts [?a, ?b] 2 tl vts ?l ?p-tl length vts]* **by** *auto*
moreover have $\text{path-image } ?p' \subseteq \text{path-image } ?p\text{-tl}$
using *sublist-path-image-subset*
by (*metis add.commute ij le-add2 length-drop length-take less-diff-conv min.absorb4 min.cobounded1 min-def vts'-sublist-tl*)

ultimately show *?thesis* **by** *blast*
qed
ultimately have $*$: $z = ?a \vee z = ?b$ **by** *blast*

let $?i = ?i'$
let $?j = ?j' - ?i' + 1$
let $?k = ?i + ?j$
let $?x1 = (2^{?i} - 1) / (2^{?i}) :: \text{real}$
let $?x2 = (2^{?(?k-1)} - 1) / (2^{?(?k-1)}) :: \text{real}$

have $?vts' = \text{take } ?j \text{ (drop } ?i \text{ vts)}$ **by** *blast*
moreover have $?k \leq \text{length } vts - 1 \wedge 2 \leq ?j$ **using** *ij* **by** *linarith*
ultimately have *path-image* $?p' = p\{?x1..?x2\}$
using *vts-sublist-path-image* *assms(1)* **unfolding** *polygon-of-def* **by** *metis*
moreover have $x1x2$: $?x1 > 1/2 \wedge ?x2 < 1$
proof-
have $?i' \geq 2$ **using** *ij* **by** *linarith*
then have $(1 :: \text{real}) < 2^{?i'} - 1$
by (*smt (z3) dual-order.strict-trans1 linorder-le-less-linear numeral-le-one-iff*
power-one-right power-strict-increasing semiring-norm(69))
thus *?thesis* **by** *simp*
qed
moreover have $p\ 0 \notin p\{?x1..?x2\} \wedge p\ (1/2) \notin p\{?x1..?x2\}$
proof-
have *False* **if** $*$: $p\ 0 \in p\{?x1..?x2\}$
proof-
obtain t **where** $t: t \in \{?x1..?x2\} \wedge p\ t = p\ 0$ **using** $*$ **by** *auto*
then have $t \geq ?x1 \wedge t \leq ?x2$ **by** *presburger*
then have $1/2 < t \wedge t < 1$ **using** $x1x2$ **by** *arg0*
thus *False*
using t *assms(1)* **unfolding** *polygon-of-def* *polygon-def* *simple-path-def*
loop-free-def
by *force*
qed
moreover have *False* **if** $*$: $p\ (1/2) \in p\{?x1..?x2\}$
proof-
obtain t **where** $t: t \in \{?x1..?x2\} \wedge p\ t = p\ (1/2)$ **using** $*$ **by** *auto*
then have $t \geq ?x1 \wedge t \leq ?x2$ **by** *presburger*
then have $1/2 < t \wedge t < 1$ **using** $x1x2$ **by** *arg0*
thus *False*
using t *assms(1)* **unfolding** *polygon-of-def* *polygon-def* *simple-path-def*
loop-free-def
by *fastforce*
qed
ultimately show *?thesis* **by** *fast*
qed
moreover have $?a = p\ 0$
by (*metis* *assms(1)* *card.empty* *empty-set* *not-numeral-le-zero* *pathstart-def*
polygon-at-least-3-vertices *polygon-of-def* *polygon-pathstart*)

moreover have $?b = p (1/2)$
proof–
have $p = ?l +++ (make\text{-}polygonal\text{-}path (tl\ vts))$
by (*smt* (*verit*, *best*) *One-nat-def Suc-1* *assms(1)* *ij length-Cons length-greater-0-conv*
length-tl less-imp-le-nat list.sel(3) list.size(3) make-polygonal-path.elims nth-Cons-0
nth-tl order-less-le-trans polygon-of-def pos2 zero-less-diff)
then have $p (1/2) = ?l 1$
unfolding *joinpaths-def* **by** *simp*
thus $?thesis$ **by** (*simp add: linepath-1'*)
qed
ultimately have $?a \notin path\text{-}image\ ?p' \wedge ?b \notin path\text{-}image\ ?p'$ **by** *presburger*
thus *False* **using** $z *$ **by** *blast*
qed
then have $frontier\ ?P1 \cap ?H \subseteq frontier\ ?H \vee frontier\ ?P2 \cap ?H \subseteq frontier\ ?H$
using *frontier-int-subset* **by** *auto*
moreover have $?L \subseteq frontier\ ?P1 \wedge ?L \subseteq frontier\ ?P2$
using *frontier-P1 frontier-P2 L-subset-F* **by** *presburger*
ultimately show $?thesis$ **using** *L-subset-H* **by** *fast*
qed

lemma *vts-on-convex-frontier-aux'*:

assumes *polygon-of p vts*
assumes $set\ vts \subseteq frontier\ (convex\ hull\ (set\ vts))$
shows $path\text{-}image\ (linepath\ (vts!0)\ (vts!1)) \subseteq frontier\ (convex\ hull\ (set\ vts))$
proof–
let $?a = vts!0$
let $?f = \lambda v. v + (-?a)$
let $?vts' = map\ ?f\ vts$
let $?p' = make\text{-}polygonal\text{-}path\ ?vts'$

have *len-vts: length vts ≥ 2*
using *assms(1) polygon-of-def polygon-vertices-length-at-least-4* **by** *fastforce*
then have $p': ?p' = ?f \circ p$
using *make-polygonal-path-translate[of vts - ?a] assms* **unfolding** *polygon-of-def*
by *presburger*
then have $0: ?vts!0 = 0$
by (*metis len-vts neg-eq-iff-add-eq-0 nth-map order-less-le-trans pos2*)
moreover have $vts': set\ ?vts' = ?f\ '(set\ vts)$ **by** *simp*
ultimately have $convex\ hull\ (set\ ?vts') = ?f\ '(convex\ hull\ (set\ vts))$
using *convex-hull-translation[of -?a set vts]* **by** *force*
then have $frontier\ (convex\ hull\ (set\ ?vts')) = frontier\ (?f\ '(convex\ hull\ (set\ vts)))$
by *auto*
then have *frontier-translation:*
 $frontier\ (convex\ hull\ (set\ ?vts')) = ?f\ '(frontier\ ((convex\ hull\ (set\ vts))))$
using *frontier-translation[of -?a convex hull (set vts)]* **by** *simp*

have $?f\ (vts!0) = ?vts!0 \wedge ?f\ (vts!1) = ?vts!1$ **using** $0\ len\text{-}vts$ **by** *auto*
then have *linepath-translation:*

$?f \text{ ' path-image (linepath (vts!0) (vts!1)) = path-image (linepath (?vts!0) (?vts!1))$
using *linepath-translation*[of $?a - ?a \text{ vts!1}$] **by** (*simp add: path-image-compose*)
have *polygon-of* $?p' \text{ ?vts'}$ **using** *translation-is-polygon* *assms(1)* p' **by** *presburger*
moreover **have** $\text{set } ?vts' \subseteq \text{frontier (convex hull (set ?vts'))}$
proof–
have $\text{frontier (convex hull (set ?vts')) = frontier (convex hull (?f \text{ ' (set vts)))}$
using vts' **by** *presburger*
then **have** $\text{frontier (convex hull (set ?vts')) = ?f \text{ ' (frontier (convex hull (set vts)))}$
using *frontier-translation* **by** *presburger*
thus $?thesis$ **using** vts' *assms(2)* **by** *auto*
qed
ultimately **have** $\text{path-image (linepath (?vts!0) (?vts!1))} \subseteq \text{frontier (convex hull (set ?vts'))}$
using *vts-on-convex-frontier-aux* *assms 0* **by** *blast*
then **have** $?f \text{ ' path-image (linepath (vts!0) (vts!1))} \subseteq ?f \text{ ' (frontier ((convex hull (set vts))))}$
using *linepath-translation* *frontier-translation* **by** *argo*
thus $?thesis$ **by** *force*
qed

lemma *vts-on-convex-frontier*:

assumes *polygon-of* $p \text{ vts}$
assumes $\text{set } vts \subseteq \text{frontier (convex hull (set vts))}$
assumes $i < \text{length } vts - 1$
shows $\text{path-image (linepath (vts!i) (vts!(i+1)))} \subseteq \text{frontier (convex hull (set vts))}$
proof–
let $?vts' = \text{rotate-polygon-vertices } vts \text{ } i$
let $?p' = \text{make-polygonal-path } ?vts'$
have *polygon-of* $?p' \text{ ?vts'}$
using *assms(1)* *polygon-of-def* *rotation-is-polygon* **by** *blast*
moreover **have** $\text{set } ?vts' \subseteq \text{frontier (convex hull (set ?vts'))}$
using *assms(1)* *assms(2)* *polygon-of-def* *rotate-polygon-vertices-same-set* **by** *auto*
ultimately **have** $\text{path-image (linepath (?vts!0) (?vts!1))} \subseteq \text{frontier (convex hull (set ?vts'))}$
using *vts-on-convex-frontier-aux'* **by** *presburger*
moreover **have** $?vts!0 = vts!i \wedge ?vts!1 = vts!(i+1)$
using *assms(3)*
using *rotated-polygon-vertices*[of $?vts' \text{ vts } i \text{ } i+1$]
using *rotated-polygon-vertices*[of $?vts' \text{ vts } i \text{ } i$]
by (*smt (verit, best) Suc-leI add.commute add.right-neutral add-2-eq-Suc'*
add-diff-cancel-left' add-lessD1 assms(1) have-wraparound-vertex hd-Nil-eq-last hd-conv-nth
last-snoc le-add1 less-diff-conv plus-1-eq-Suc polygon-of-def)
moreover **have** $\text{frontier (convex hull (set ?vts'))} = \text{frontier (convex hull (set vts))}$

by (metis assms(1) polygon-of-def rotate-polygon-vertices-same-set)
ultimately show ?thesis by argo
qed

lemma *pts-on-frontier-means-path-image-on-frontier*:
assumes *polygon-of* p *pts*
assumes $set\ pts \subseteq frontier\ (convex\ hull\ (set\ pts))$
shows $path\ image\ p \subseteq frontier\ (convex\ hull\ (set\ pts))$
proof(rule *subsetI*)
let $?H = convex\ hull\ (set\ pts)$
fix x **assume** $x \in path\ image\ p$
moreover have $path\ image\ p = (\bigcup \{path\ image\ (linepath\ (pts!i)\ (pts!(i+1))) \mid$
 $i. i \leq (length\ pts) - 2\})$
using *polygonal-path-image-linepath-union* *assms* **unfolding** *polygon-of-def*
by (metis (no-types, lifting) *add-leD2 numeral-Bit0 polygon-vertices-length-at-least-4*)
ultimately obtain i **where** $i \leq (length\ pts) - 2 \wedge x \in path\ image\ (linepath$
 $(pts!i)\ (pts!(i+1)))$
by *blast*
thus $x \in frontier\ ?H$
by (smt (verit, ccfv-SIG) *One-nat-def Suc-diff-Suc add.commute add-2-eq-Suc'*
assms(1) assms(2) in-mono le-add1 le-zero-eq less-Suc-eq-le less-diff-conv linorder-not-less
plus-1-eq-Suc pts-on-convex-frontier pts-on-convex-frontier-aux')
qed

lemma *pts-on-convex-frontier-interior*:
assumes *polygon-of* p *pts*
assumes $set\ pts \subseteq frontier\ (convex\ hull\ (set\ pts))$
shows $path\ inside\ p = interior\ (convex\ hull\ (set\ pts))$
proof–
let $?H = convex\ hull\ (set\ pts)$

have $path\ inside\ p \subseteq interior\ (convex\ hull\ (set\ pts))$
by (metis (no-types, lifting) *Un-empty assms(1) convex-contains-simple-closed-path-imp-contains-path-inside*
convex-convex-hull convex-hull-eq-empty convex-hull-of-polygon-is-convex-hull-of-pts
empty-set inside-outside-def inside-outside-polygon interior-maximal length-greater-0-conv
polygon-def polygon-of-def polygon-path-image-subset-convex)
moreover have $interior\ (convex\ hull\ (set\ pts)) \subseteq path\ inside\ p$
proof(rule *ccontr*)
assume $*$: $\neg interior\ (convex\ hull\ (set\ pts)) \subseteq path\ inside\ p$
then obtain x **where** $x \in interior\ (convex\ hull\ (set\ pts)) - path\ inside\ p$
by *blast*
obtain y **where** $y \in path\ inside\ p$
using *inside-outside-polygon* *assms* **unfolding** *inside-outside-def polygon-of-def*
by *fastforce*

let $?l = linepath\ x\ y$
have $1: path\ image\ ?l \subseteq interior\ ?H$
by (metis (no-types, lifting) *DiffE calculation convex-contains-segment con-*
vox-convex-hull convex-interior in-mono linepath-image-01 path-defs(4) x y)

have $\text{path-image } ?l \cap \text{frontier } (\text{path-inside } p) \neq \{\}$
using *inside-outside-polygon* *assms* **unfolding** *inside-outside-def* *polygon-of-def*
by (*smt* (*verit*) * *Diff-disjoint* *Diff-eq-empty-iff* *Int-Un-eq*(2) *Int-assoc* *Un-Int-eq*(3)
assms(1) *calculation* *connected-Int-frontier* *convex-connected* *convex-convex-hull* *convex-interior* *frontier-def* *inf.absorb-iff2* *vts-on-frontier-means-path-image-on-frontier*)
then have 2: $\text{path-image } ?l \cap \text{path-image } p \neq \{\}$
using *inside-outside-polygon* *assms* **unfolding** *inside-outside-def* *polygon-of-def*
by *blast*

show *False*
using 1 2 *vts-on-frontier-means-path-image-on-frontier*
using *Diff-disjoint* *Int-lower2* *Int-subset-iff* *assms*(1) *assms*(2) *frontier-def*
inf-le1
by *fastforce*
qed
ultimately show *?thesis* **by** *blast*
qed

lemma *vts-subset-frontier*:
assumes *polygon-of* *p* *vts*
assumes $\text{set } vts \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$
shows $\text{convex } (\text{path-image } p \cup \text{path-inside } p)$
by (*metis* *assms*(1) *assms*(2) *vts-on-convex-frontier-interior* *convex-convex-hull* *convex-interior* *polygon-convex-iff* *polygon-of-def* *sup-commute*)

lemma *convex-hull-of-nonconvex-polygon-strict-subset-ep*:
assumes *polygon-of* *p* *vts*
assumes $\neg (\text{convex } (\text{path-image } p \cup \text{path-inside } p))$
shows $\{v. v \text{ extreme-point-of } (\text{convex hull } (\text{set } vts))\} \subset \text{set } vts$
proof –
let *?ep* = $\{v. v \text{ extreme-point-of } (\text{convex hull } (\text{set } vts))\}$
let *?H* = $\text{convex hull } (\text{set } vts)$
have *?ep* $\subseteq \text{frontier } ?H$
by (*metis* *Krein-Milman-frontier* *List.finite-set* *convex-convex-hull* *extreme-point-of-convex-hull* *finite-imp-compact-convex-hull* *mem-Collect-eq* *subsetI*)
thus *?thesis* **using** *assms* *vts-subset-frontier* *extreme-points-of-convex-hull* **by**
force
qed

lemma *convex-hull-of-nonconvex-polygon-strict-subset*:
assumes *polygon-of* *p* *vts*
assumes $\neg (\text{convex } (\text{path-image } p \cup \text{path-inside } p))$
shows $\exists v \in \text{set } vts. v \in \text{interior } (\text{convex hull } (\text{set } vts))$
using *assms* *vts-subset-frontier*
by (*smt* (*verit*) *Diff-iff* *UnCI* *closure-Un-frontier* *frontier-def* *hull-inc* *subsetI*)

lemma *convex-polygon-means-linepaths-inside*:
fixes *p* :: *R-to-R2*
assumes *polygon-of* *p* *vts*

```

assumes convex-is:  $\text{convex hull (set vts)} = (\text{path-inside } p \cup \text{path-image } p)$ 
assumes a-in:  $a \in (\text{path-inside } p \cup \text{path-image } p)$ 
assumes b-in:  $b \in (\text{path-inside } p \cup \text{path-image } p)$ 
shows path-image (linepath a b)  $\subseteq (\text{path-inside } p \cup \text{path-image } p)$ 
proof –
  let ?conv =  $\text{path-inside } p \cup \text{path-image } p$ 
  have  $\forall u \geq 0. \forall v \geq 0. u + v = 1 \longrightarrow u *_R a + v *_R b \in ?conv$ 
    using convex-is a-in b-in unfolding convex-def
    by (metis (no-types, lifting) convexD convex-convex-hull convex-is)
  then have  $(1 - x) *_R a + x *_R b \in ?conv$  if x-in:  $x \in \{0..1\}$  for x
    using x-in by auto
  then show thesis unfolding linepath-def path-image-def
    by fast
qed

```

```

end
theory Polygon-Splitting
imports
  HOL-Analysis.Complete-Measure
  Polygon-Jordan-Curve
  Polygon-Convex-Lemmas
begin

```

19 Polygon Splitting

```

lemma split-up-a-list-into-3-parts:
  fixes i j:: nat
  assumes  $i < \text{length vts} \wedge j < \text{length vts} \wedge i < j$ 
  shows
     $\text{vts} = (\text{take } i \text{ vts}) @ ((\text{vts} ! i) \# ((\text{take } (j - i - 1) (\text{drop } (\text{Suc } i) \text{ vts})) @ (\text{vts} ! j) \# \text{drop } (j - i) (\text{drop } (\text{Suc } i) \text{ vts})))$ 
proof –
  let ?x =  $\text{vts} ! i$ 
  let ?y =  $\text{vts} ! j$ 
  let ?vts1 =  $\text{take } i \text{ vts}$ 
  let ?drop-list =  $\text{drop } (\text{Suc } i) \text{ vts}$ 
  have vts-is:  $\text{vts} = ?vts1 @ \text{vts} ! i \# \text{drop } (\text{Suc } i) \text{ vts}$ 
    using split-list assms
    by (meson id-take-nth-drop)
  then have len-vts1:  $\text{length } ?vts1 = i$ 
    using length-take[of i vts] assms
    by auto
  have gt-eq:  $j - i - 1 \geq 0$ 
    using assms by auto
  let ?ind =  $j - i - 1$ 
  have drop-is:  $\text{drop } (\text{Suc } i) \text{ vts} ! (j - i - 1) = ?y$ 
    using assms by auto
  then have drop-list-is:  $?drop-list = \text{take } ?ind ?drop-list @ ?y \# (\text{drop } (j - i) ?drop-list)$ 

```

by (*metis Suc-diff-Suc Suc-leI assms diff-Suc-1 diff-less-mono id-take-nth-drop length-drop*)
have $\text{length } (\text{drop } (\text{Suc } ?\text{ind}) ?\text{drop-list}) = \text{length } \text{vts} - j - 1$
using $\text{length-drop}[\text{of } \text{Suc } (j - i - 1) (\text{drop } (\text{Suc } i) \text{vts})]$ *length-take assms*
by *auto*
then show *?thesis*
using *vts-is drop-list-is len-vts1*
by *presburger*
qed

definition *is-polygon-cut* :: (real^2) list $\Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{bool}$ **where**
is-polygon-cut vts x y =
 $(x \neq y \wedge$
 $\text{polygon } (\text{make-polygonal-path } \text{vts}) \wedge$
 $\{x, y\} \subseteq \text{set } \text{vts} \wedge$
 $\text{path-image } (\text{linepath } x \ y) \cap \text{path-image } (\text{make-polygonal-path } \text{vts}) = \{x, y\} \wedge$
 $\text{path-image } (\text{linepath } x \ y) \cap \text{path-inside } (\text{make-polygonal-path } \text{vts}) \neq \{\})$

definition *is-polygon-cut-path* :: (real^2) list $\Rightarrow R\text{-to-}R^2 \Rightarrow \text{bool}$ **where**
is-polygon-cut-path vts cutpath =
 $(\text{let } x = \text{pathstart } \text{cutpath} ; y = \text{pathfinish } \text{cutpath} \text{ in}$
 $(x \neq y \wedge$
 $\text{polygon } (\text{make-polygonal-path } \text{vts}) \wedge$
 $\{x, y\} \subseteq \text{set } \text{vts} \wedge$
 $\text{simple-path } \text{cutpath} \wedge$
 $\text{path-image } \text{cutpath} \cap \text{path-image } (\text{make-polygonal-path } \text{vts}) = \{x, y\} \wedge$
 $\text{path-image } \text{cutpath} \cap \text{path-inside } (\text{make-polygonal-path } \text{vts}) \neq \{\})$

definition *is-polygon-split* ::
 (real^2) list $\Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{bool}$ **where**
is-polygon-split vts i j =
 $(i < \text{length } \text{vts} \wedge j < \text{length } \text{vts} \wedge i < j \wedge$
 $(\text{let } \text{vts1} = (\text{take } i \ \text{vts}) \text{ in}$
 $\text{let } \text{vts2} = (\text{take } (j - i - 1) (\text{drop } (\text{Suc } i) \ \text{vts})) \text{ in}$
 $\text{let } \text{vts3} = \text{drop } (j - i) (\text{drop } (\text{Suc } i) \ \text{vts}) \text{ in}$
 $\text{let } x = \text{vts} ! i \text{ in}$
 $\text{let } y = \text{vts} ! j \text{ in}$
 $\text{let } p = \text{make-polygonal-path } (\text{vts}@[\text{vts}!0]) \text{ in}$
 $\text{let } p1 = \text{make-polygonal-path } (x\#(\text{vts2}@[y, x])) \text{ in}$
 $\text{let } p2 = \text{make-polygonal-path } (\text{vts1} @ [x, y] @ \text{vts3} @ [\text{vts} ! 0]) \text{ in}$
 $\text{let } c1 = \text{make-polygonal-path } (x\#(\text{vts2}@[y])) \text{ in}$
 $\text{let } c2 = \text{make-polygonal-path } (\text{vts1} @ [x, y] @ \text{vts3}) \text{ in}$
 $(\text{is-polygon-cut } (\text{vts}@[\text{vts}!0]) \ x \ y \wedge$
 $\text{polygon } p \wedge \text{polygon } p1 \wedge \text{polygon } p2 \wedge$
 $\text{path-inside } p1 \cap \text{path-inside } p2 = \{\} \wedge$
 $\text{path-inside } p1 \cup \text{path-inside } p2 \cup (\text{path-image } (\text{linepath } x \ y) - \{x, y\}) =$
 $\text{path-inside } p$

$$\begin{aligned}
& \wedge ((\text{path-image } p1) - (\text{path-image } (\text{linepath } x \ y))) \cap ((\text{path-image } p2) - \\
& (\text{path-image } (\text{linepath } x \ y))) \\
& = \{\} \\
& \wedge \text{path-image } p \\
& = ((\text{path-image } p1) - (\text{path-image } (\text{linepath } x \ y))) \cup ((\text{path-image } p2) - \\
& (\text{path-image } (\text{linepath } x \ y))) \cup \{x, y\} \\
&)))
\end{aligned}$$

definition *is-polygon-split-path* :: (real²) list ⇒ nat ⇒ nat ⇒ (real²) list ⇒ bool where

$$\begin{aligned}
& \text{is-polygon-split-path } vts \ i \ j \ \text{cutvts} = \\
& (i < \text{length } vts \wedge j < \text{length } vts \wedge i < j \wedge \\
& (\text{let } vts1 = (\text{take } i \ vts) \ \text{in} \\
& \text{let } vts2 = (\text{take } (j - i - 1) \ (\text{drop } (\text{Suc } i) \ vts)) \ \text{in} \\
& \text{let } vts3 = \text{drop } (j - i) \ (\text{drop } (\text{Suc } i) \ vts) \ \text{in} \\
& \text{let } x = vts!i \ \text{in} \\
& \text{let } y = vts!j \ \text{in} \\
& \text{let } \text{cutpath} = \text{make-polygonal-path } (x \# \ \text{cutvts} \ @ \ [y]) \ \text{in} \\
& \text{let } p = \text{make-polygonal-path } (vts@[vts!0]) \ \text{in} \\
& \text{let } p1 = \text{make-polygonal-path } (x\#(vts2 \ @ \ [y] \ @ \ (\text{rev } \text{cutvts}) \ @ \ [x])) \ \text{in} \\
& \text{let } p2 = \text{make-polygonal-path } (vts1 \ @ \ ([x] \ @ \ \text{cutvts} \ @ \ [y]) \ @ \ vts3 \ @ \ [vts!0]) \ \text{in} \\
& \text{let } c1 = \text{make-polygonal-path } (x\#(vts2@[y])) \ \text{in} \\
& \text{let } c2 = \text{make-polygonal-path } (vts1 \ @ \ ([x] \ @ \ \text{cutvts} \ @ \ [y]) \ @ \ vts3) \ \text{in} \\
& (\text{is-polygon-cut-path } (vts@[vts!0]) \ \text{cutpath} \ \wedge \\
& \text{polygon } p \ \wedge \text{polygon } p1 \ \wedge \text{polygon } p2 \ \wedge \\
& \text{path-inside } p1 \ \cap \ \text{path-inside } p2 = \{\} \ \wedge \\
& \text{path-inside } p1 \ \cup \ \text{path-inside } p2 \ \cup \ (\text{path-image } \text{cutpath} - \{x, y\}) = \text{path-inside} \\
& p \\
& \wedge ((\text{path-image } p1) - (\text{path-image } \text{cutpath})) \cap ((\text{path-image } p2) - (\text{path-image} \\
& \text{cutpath})) = \{\} \\
& \wedge \text{path-image } p \\
& = ((\text{path-image } p1) - (\text{path-image } \text{cutpath})) \cup ((\text{path-image } p2) - (\text{path-image} \\
& \text{cutpath})) \cup \{x, y\} \\
&)))
\end{aligned}$$

lemma *polygon-split-add-measure*:

fixes $p \ p1 \ p2$:: R-to-R2

assumes *is-polygon-split* $vts \ i \ j$

assumes $vts1 = (\text{take } i \ vts)$

$vts2 = (\text{take } (j - i - 1) \ (\text{drop } (\text{Suc } i) \ vts))$

$vts3 = \text{drop } (j - i) \ (\text{drop } (\text{Suc } i) \ vts)$

$x = vts!i$

$y = vts!j$

$p = \text{make-polygonal-path } (vts@[vts!0])$

$p1 = \text{make-polygonal-path } (x\#(vts2@[y, x]))$

$p2 = \text{make-polygonal-path } (vts1 \ @ \ [x, y] \ @ \ vts3 \ @ \ [vts!0])$

defines $M1 \equiv \text{measure lebesgue } (\text{path-inside } p1)$ **and**

$M2 \equiv \text{measure lebesgue } (\text{path-inside } p2)$ **and**

$M \equiv \text{measure lebesgue } (\text{path-inside } p)$

shows $M1 + M2 = M$
proof –
let $?cut = \text{linepath } x \ y$
let $?cut\text{-open-image} = (\text{path-image } ?cut) - \{x, y\}$
let $?P = \text{path-inside } p$
let $?P1 = \text{path-inside } p1$
let $?P2 = \text{path-inside } p2$
let $?M = \text{space lebesgue}$
let $?A = \text{sets lebesgue}$
let $?μ = \text{emeasure lebesgue}$

have open $?P1$
by (*metis* *assms*(1) *assms*(3) *assms*(5) *assms*(6) *assms*(8) *closed-path-image is-polygon-split-def open-inside path-inside-def polygon-def simple-path-def*)
then have $P1\text{-measurable: } ?P1 \in ?A$ **by** *simp*

have open $?P2$
by (*metis* *assms*(1) *assms*(2) *assms*(4) *assms*(5) *assms*(6) *assms*(9) *closed-path-image is-polygon-split-def open-inside path-inside-def polygon-def simple-path-def*)
then have $P2\text{-measurable: } ?P2 \in ?A$ **by** *simp*

have $?P1 \cap ?P2 = \{\}$
by (*metis* *assms*(1) *assms*(2) *assms*(3) *assms*(4) *assms*(5) *assms*(6) *assms*(8) *assms*(9) *is-polygon-split-def*)
then have *sum-union-finite:* $?μ ?P1 + ?μ ?P2 = ?μ (?P1 \cup ?P2)$
using *plus-emeasure P1-measurable P2-measurable* **by** *blast*

have *measure lebesgue* $?P1 = ?μ ?P1$
by (*metis* *assms*(1) *assms*(3) *assms*(5) *assms*(6) *assms*(8) *bounded-inside bounded-set-imp-lmeasurable bounded-simple-path-image emeasure-eq-ennreal-measure emeasure-notin-sets ennreal-0 fmeasurableD2 is-polygon-split-def measure-zero-top path-inside-def polygon-def*)
moreover have *measure lebesgue* $?P2 = ?μ ?P2$
by (*metis* *Sigma-Algebra.measure-def* *assms*(1) *assms*(2) *assms*(4) *assms*(5) *assms*(6) *assms*(9) *bounded-inside bounded-path-image bounded-set-imp-lmeasurable emeasure-eq-ennreal-measure emeasure-notin-sets enn2real-top ennreal-0 fmeasurableD2 is-polygon-split-def path-inside-def polygon-def simple-path-def*)
ultimately have $?μ (?P1 \cup ?P2) = M1 + M2$
using *assms*(10) *assms*(11) *sum-union-finite* **by** *auto*
moreover have $?μ (?P1 \cup ?P2) = ?μ ?P$

proof –
have $?μ (\text{path-image } ?cut) = 0$ **using** *linepath-has-emeasure-0* **by** *blast*
then have $(\text{path-image } ?cut) \in \text{null-sets lebesgue}$ **by** *auto*
moreover have $\{x, y\} \in \text{null-sets lebesgue}$ **by** *simp*
ultimately have $?cut\text{-open-image} \in \text{null-sets lebesgue}$ **using** *measure-Diff-null-set*
by *auto*
moreover have $?P = ?P1 \cup ?P2 \cup ?cut\text{-open-image}$
by (*metis* *assms*(1) *assms*(2) *assms*(3) *assms*(4) *assms*(5) *assms*(6) *assms*(7) *assms*(8) *assms*(9) *is-polygon-split-def*)

```

ultimately show ?thesis
  by (simp add: P1-measurable P2-measurable emeasure-Un-null-set sets.Un)
qed
ultimately show ?thesis
  by (smt (verit, best) M1-def M2-def M-def emeasure-eq-ennreal-measure enn2real-ennreal
ennreal-neq-top measure-nonneg)
qed

lemma polygonal-paths-measurable:
  shows path-image (make-polygonal-path vts) ∈ sets lebesgue
proof (induct vts rule: make-polygonal-path-induct)
  case (Empty ell)
  then show ?case by auto
next
  case (Single ell)
  then obtain a where ell = [a]
  by (metis Cons-nth-drop-Suc One-nat-def drop0 drop-eq-Nil le-numeral-extra(4)
zero-less-one)
  then show ?case using make-polygonal-path.simps(2)[of a] by simp
next
  case (Two ell)
  then obtain a b where ell = [a, b]
  by (metis Cons-nth-drop-Suc One-nat-def Suc-1 append-Nil drop-eq-Nil2 dual-order.refl
id-take-nth-drop lessI pos2 take0)
  then show ?case using make-polygonal-path.simps(3)[of a b] by simp
next
  case (Multiple ell)
  then have ell = (ell ! 0) # (ell ! 1) # (ell ! 2) # (drop 3 ell)
  by (metis Cons-nth-drop-Suc One-nat-def Suc-1 drop0 le-Suc-eq linorder-not-less
numeral-3-eq-3)
  then have make-polygonal-path ell =
    linepath (ell ! 0) (ell ! 1) +++ make-polygonal-path (ell ! 1 # ell ! 2 # (drop
3 ell))
  by (metis make-polygonal-path.simps(4))

  then have path-image (make-polygonal-path ell) = path-image (linepath (ell ! 0)
(ell ! 1)) ∪ path-image (make-polygonal-path (ell ! 1 # ell ! 2 # (drop 2 ell)))
  using Cons-nth-drop-Suc Multiple.hyps(1) One-nat-def Suc-1 Un-assoc ⟨ell =
ell ! 0 # ell ! 1 # ell ! 2 # drop 3 ell⟩ list.discI make-polygonal-path.simps(2)
make-polygonal-path.simps(3) nth-Cons-0 numeral-3-eq-3 path-image-cons-union
proof -
  have f1: ell = ell ! 0 # ell ! 1 # ell ! Suc 1 # drop 3 ell
  using Suc-1 ⟨ell = ell ! 0 # ell ! 1 # ell ! 2 # drop 3 ell⟩ by presburger
  have Suc 1 < length ell
  by (smt (z3) Suc-1 ⟨2 < length ell⟩)
  then have f2: drop (Suc 1) ell = ell ! Suc 1 # drop (Suc (Suc 1)) ell
  by (smt (z3) Cons-nth-drop-Suc)
  have f3: ∀ v va vs. path-image (make-polygonal-path (v # va # vs)) = path-image
(linepath v va) ∪ path-image (make-polygonal-path (va # vs))

```

```

    by (metis (no-types) list.discI nth-Cons-0 path-image-cons-union)
    have f4:  $\forall V v va. \text{path-image } (\text{linepath } (v::(\text{real}, 2) \text{ vec}) va) \cup (\text{path-image } (\text{linepath } va va) \cup V) = \text{path-image } (\text{linepath } v va) \cup V$ 
    by auto
    have path-image (make-polygonal-path ell) = path-image (make-polygonal-path (ell ! 0 # ell ! 1 # drop (Suc 1) ell))
    using f2 f1 by (simp add: numeral-3-eq-3)
    then have path-image (make-polygonal-path ell) = path-image (linepath (ell ! 0) (ell ! 1))  $\cup$  path-image (make-polygonal-path (ell ! 1 # ell ! Suc 1 # drop (Suc 1) ell))
    using f4 f3 f2 by presburger
    then show ?thesis
    using Suc-1 by presburger
qed
then show ?case using Multiple(3)
by (metis (no-types, lifting) Cons-nth-drop-Suc Multiple.hyps(1) Multiple.hyps(2) One-nat-def Suc-1  $\langle$ ell = ell ! 0 # ell ! 1 # ell ! 2 # drop 3 ell $\rangle$  list.discI make-polygonal-path.simps(3) nth-Cons-0 numeral-3-eq-3 path-image-cons-union sets.Un)

```

qed

lemma *polygonal-path-has-emeasure-0:*

```

    shows emeasure lebesgue (path-image (make-polygonal-path vts)) = 0
proof (induct vts)
  case Nil
    then show ?case by auto
next
  case (Cons a vts)
    then show ?case
    by (metis linepath-is-negligible make-polygonal-path.simps(2) negligible-Un negligible-iff-emeasure0 path-image-cons-union polygonal-paths-measurable)
qed

```

lemma *polygon-split-path-add-measure:*

```

fixes p p1 p2 :: R-to-R2
assumes is-polygon-split-path vts i j cutvts
assumes vts1 = (take i vts)
          vts2 = (take (j - i - 1) (drop (Suc i) vts))
          vts3 = drop (j - i) (drop (Suc i) vts)
          x = vts ! i
          y = vts ! j
          p = make-polygonal-path (vts@[vts!0])
          p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x]))
          p2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @ [vts ! 0])
defines M1  $\equiv$  measure lebesgue (path-inside p1) and
          M2  $\equiv$  measure lebesgue (path-inside p2) and
          M  $\equiv$  measure lebesgue (path-inside p)
shows M1 + M2 = M
proof -

```

```

let ?cut = make-polygonal-path (x # cutvts @ [y])
let ?cut-open-image = (path-image ?cut) - {x, y}
let ?P = path-inside p
let ?P1 = path-inside p1
let ?P2 = path-inside p2
let ?M = space lebesgue
let ?A = sets lebesgue
let ?μ = emeasure lebesgue

have open ?P1
  by (metis assms(1) assms(3) assms(5) assms(6) assms(8) closed-path-image
  is-polygon-split-path-def open-inside path-inside-def polygon-def simple-path-def)
  then have P1-measurable: ?P1 ∈ ?A by simp

have open ?P2
  by (metis assms(1) assms(2) assms(4) assms(5) assms(6) assms(9) closed-path-image
  is-polygon-split-path-def open-inside path-inside-def polygon-def simple-path-def)
  then have P2-measurable: ?P2 ∈ ?A by simp

have ?P1 ∩ ?P2 = {}
  by (metis assms(1) assms(2) assms(3) assms(4) assms(5) assms(6) assms(8)
  assms(9) is-polygon-split-path-def)
  then have sum-union-finite: ?μ ?P1 + ?μ ?P2 = ?μ (?P1 ∪ ?P2)
  using plus-emeasure P1-measurable P2-measurable by blast

have ?μ (path-image q) = 0 ⇒ (path-image q) ∈ null-sets lebesgue if *:
  path-image q ∈ sets lebesgue for q::real ⇒ (real, 2) vec
  using null-sets-def * by blast

have measure lebesgue ?P1 = ?μ ?P1
  by (metis Sigma-Algebra.measure-def assms(1) assms(3) assms(5) assms(6) assms(8) bounded-inside
  bounded-set-imp-lmeasurable bounded-simple-path-image emeasure-eq-ennreal-measure
  emeasure-notin-sets ennreal-0 fmeasurableD2 is-polygon-split-path-def measure-zero-top
  path-inside-def polygon-def)
  moreover have measure lebesgue ?P2 = ?μ ?P2
  by (metis Sigma-Algebra.measure-def assms(1) assms(2) assms(4) assms(5)
  assms(6) assms(9) bounded-inside bounded-path-image bounded-set-imp-lmeasurable
  emeasure-eq-ennreal-measure emeasure-notin-sets enn2real-top ennreal-0 fmeasur-
  ableD2 is-polygon-split-path-def path-inside-def polygon-def simple-path-def)
  ultimately have ?μ (?P1 ∪ ?P2) = M1 + M2
  using assms(10) assms(11) sum-union-finite by auto
  moreover have ?μ (?P1 ∪ ?P2) = ?μ ?P
  proof –
    have ?μ (path-image ?cut) = 0 using polygonal-path-has-emeasure-0
    by presburger
    then have (path-image ?cut) ∈ null-sets lebesgue using polygonal-paths-measurable
    by blast
    moreover have {x, y} ∈ null-sets lebesgue by simp
    ultimately have ?cut-open-image ∈ null-sets lebesgue using measure-Diff-null-set

```

```

by auto
  moreover have ?P = ?P1 ∪ ?P2 ∪ ?cut-open-image
  by (metis assms(1) assms(2) assms(3) assms(4) assms(5) assms(6) assms(7)
      assms(8) assms(9) is-polygon-split-path-def)
  ultimately show ?thesis
  by (simp add: P1-measurable P2-measurable emeasure-Un-null-set sets.Un)
qed
ultimately show ?thesis
by (smt (verit, best) M1-def M2-def M-def emeasure-eq-ennreal-measure enn2real-ennreal
    ennreal-neq-top measure-nonneg)
qed

```

lemma *polygon-cut-path-to-split-path-vtx0*:

```

fixes p :: R-to-R2
assumes polygon-p: polygon p and
  i-gt: i > 0 and
  i-lt: i < length vts and
  p-is: p = make-polygonal-path (vts @ [vts ! 0]) and
  cutpath: cutpath = make-polygonal-path ([vts!0] @ cutvts @ [vts!i]) and
  have-cut: is-polygon-cut-path (vts @ [vts!0]) cutpath
shows is-polygon-split-path vts 0 i cutvts
proof -
  let ?vts2 = take (i - 1) (drop 1 vts)
  let ?vts3 = drop i (drop 1 vts)
  let ?x = vts ! 0
  let ?y = vts ! i

  let ?c3-vts = [?x] @ cutvts @ [?y]
  let ?c3 = cutpath
  let ?c3-rev-vts = rev ?c3-vts
  let ?c3-rev = make-polygonal-path ?c3-rev-vts
  let ?c3' = reversepath ?c3

  let ?p = make-polygonal-path (vts @ [vts ! 0])
  let ?p1-vts = ?x # ?vts2 @ ?c3-rev-vts
  let ?p1 = make-polygonal-path ?p1-vts
  let ?p1-rot-vts = ?c3-rev-vts @ ?vts2 @ [?y]
  let ?p1-rot = make-polygonal-path ?p1-rot-vts
  let ?p2-vts = ?c3-vts @ ?vts3 @ [?x]
  let ?p2 = make-polygonal-path ?p2-vts
  let ?c1-vts = ?x # ?vts2 @ [?y]
  let ?c1 = make-polygonal-path ?c1-vts
  let ?c2-vts = [?y] @ ?vts3 @ [?x]
  let ?c2 = reversepath (make-polygonal-path ?c2-vts)
  let ?c2'-vts = [?y] @ ?vts3 @ [?x]
  let ?c2' = (make-polygonal-path (?c2'-vts))

  have distinct-vts: distinct vts
  using polygon-p p-is

```

```

using polygon-def simple-polygonal-path-vts-distinct by force
have len-vts-gteq3: length vts ≥ 3
using polygon-p p-is polygon-vertices-length-at-least-4 by fastforce

then have ?x # ?vts2 @ [?y] = take (i+1) (vts@ [vts ! 0])
by (smt (verit, ccfv-threshold) i-gt Cons-nth-drop-Suc Suc-eq-plus1 Suc-pred'
add-less-cancel-left butlast-snoc drop0 drop-drop hd-drop-conv-nth i-lt length-append-singleton
length-greater-0-conv less-imp-le-nat linorder-not-less list.size(3) plus-1-eq-Suc take-Suc-Cons
take-all-iff take-butlast take-hd-drop)
have [?y] @ ?vts3 @ [?x] = drop (i) (vts @ [vts ! 0])
using i-gt
by (metis (no-types, lifting) Cons-eq-appendI Cons-nth-drop-Suc Suc-eq-plus1
append-Nil diff-is-0-eq' drop-0 drop-append drop-drop i-lt less-imp-le-nat)

have card-gteq: card (set vts) ≥ 3
using polygon-at-least-3-vertices-wraparound polygon-p p-is
by (metis butlast-conv-take butlast-snoc)
then have vts ≠ []
by auto
then have vts-is: vts = ?x # ?vts2 @ ?y # ?vts3
using split-up-a-list-into-3-parts[of 0 vts i] i-gt i-lt
by auto

have elem-prop1: last ?c1-vts = ?y
by (metis (no-types, lifting) last.simps snoc-eq-iff-butlast)
have elem-prop2: (vts ! 0 # (rev ?vts3) @ [vts ! i]) !
(length (vts ! 0 # drop i (drop 1 vts) @ [vts ! i]) - 1) = vts ! i
by (metis diff-Suc-1 length-Cons length-append-singleton length-rev nth-Cons-Suc
nth-append-length)
have path-image cutpath = path-image ?c3' by simp
then have path-image ?p1 = path-image (?c1 +++ ?c3-rev)
using elem-prop1 assms make-polygonal-path-image-append-alt[of ?p1 ?p1-vts
?c1 ?c1-vts ?c3-rev ?c3-rev-vts]
by simp
also have ... = path-image ?c1 ∪ path-image ?c3-rev
by (metis (no-types, opaque-lifting) append-Cons append-Nil elem-prop1 hd-conv-nth
last-conv-nth list.discI list.sel(1) path-image-join polygon-pathfinish polygon-pathstart
rev.simps(2) rev-rev-ident)
finally have image-prop: path-image ?p1 = path-image ?c1 ∪ path-image cutpath
using rev-vts-path-image cutpath by presburger
have path-image ?c3' = path-image ?c3
using cutpath rev-vts-path-image by force
then have path-image-p1: path-image ?c1 ∪ path-image ?c3 = path-image ?p1
using image-prop by presburger

have ?p2-vts = ?c3-vts @ (tl ?c2-vts) by simp
then have path-image ?p2 = path-image (?c3 +++ ?c2')
using make-polygonal-path-image-append-alt[of ?p2 ?p2-vts ?c3 ?c3-vts ?c2']

```

```

?c2-vts]
  unfolding assms by auto
  then have path-image-p2: path-image ?c2  $\cup$  path-image ?c3 = path-image ?p2
  by (metis (no-types, opaque-lifting) Un-commute append-Cons append-Nil cut-
path last-conv-nth nth-Cons-0 path-image-join path-image-reversepath polygon-pathfinish
polygon-pathstart snoc-eq-iff-butlast)

  have drop 1 vts = take (i - 1) (drop 1 vts) @ [vts ! i] @ drop i (drop 1 vts)
  by (metis (no-types, lifting) Cons-eq-appendI Cons-nth-drop-Suc Suc-eq-plus1
Suc-pred' append.simps(1) append-take-drop-id drop-drop i-gt i-lt)
  then have vts-is: vts @ [vts ! 0] = vts ! 0 # take (i - 1) (drop 1 vts) @ [vts !
i] @ drop i (drop 1 vts) @ [vts ! 0]
  by (metis (no-types, opaque-lifting) Cons-nth-drop-Suc One-nat-def append.assoc
append-Cons drop0 i-lt length-pos-if-in-set nth-mem)
  let ?vts1' = take (i - 1) (drop 1 vts)
  let ?vts2' = drop i (drop 1 vts)
  have path-im-p: path-image
    (make-polygonal-path
      ((vts ! 0 # ?vts1') @ [vts ! i] @ [vts ! i] @ ?vts2' @ [vts ! 0])) =
    path-image
      (make-polygonal-path
        ((vts ! 0 # ?vts1') @ [vts ! i] @ ?vts2' @ [vts ! 0]))
  using make-polygonal-path-image-append-helper[of vts ! 0 # ?vts1' ?vts2' @
[vts ! 0]] by auto
  have path-image
    (make-polygonal-path
      ((vts ! 0 # ?vts1') @ [vts ! i] @ [vts ! i] @ ?vts2' @ [vts ! 0])) = path-image
    (make-polygonal-path ((vts ! 0 # ?vts1') @ [vts ! i]) +++ (linepath (vts ! i) (vts !
i)) +++ make-polygonal-path ([vts ! i] @ ?vts2' @ [vts ! 0]))
  using make-polygonal-path-image-append[of (vts ! 0 # ?vts1') @ [vts ! i] [vts !
i] @ ?vts2' @ [vts ! 0]]

  by (smt (verit) add-2-eq-Suc' append.assoc append-Cons diff-Suc-1 le-add2
length-Cons length-append-singleton nth-Cons-0 nth-append-length)
  then have path-image p = path-image (make-polygonal-path ((vts ! 0 # ?vts1')
@ [vts ! i]) +++ (linepath (vts ! i) (vts ! i)) +++ make-polygonal-path ([vts ! i] @
?vts2' @ [vts ! 0]))
  using path-im-p p-is vts-is
  by simp
  then have path-image p = path-image ?c1  $\cup$  path-image (linepath (vts ! i) (vts
! i))  $\cup$  path-image (make-polygonal-path ([vts ! i] @ ?vts2' @ [vts ! 0]))
  by (metis (no-types, lifting) Un-assoc append-Cons elem-prop1 list.discI nth-Cons-0
path-image-join pathfinish-linepath pathstart-join pathstart-linepath polygon-pathfinish
polygon-pathstart last-conv-nth)
  moreover have ... = path-image ?c1  $\cup$  {vts ! i}  $\cup$  path-image (make-polygonal-path
([vts ! i] @ ?vts2' @ [vts ! 0]))
  by auto
  moreover have ... = path-image ?c1  $\cup$  path-image (make-polygonal-path ([vts !
i] @ ?vts2' @ [vts ! 0]))

```

```

using vertices-on-path-image by fastforce
ultimately have path-image-p: path-image p = path-image ?c1  $\cup$  path-image
?c2
using path-image-reversepath by blast

have simple-path-polygon: simple-path (make-polygonal-path (?x # ?vts2 @ ?y
# ?vts3 @ [?x]))
using polygon-p p-is vts-is
using Cons-eq-appendI append-self-conv2 polygon-def by auto
then have loop-free-polygon: loop-free (make-polygonal-path (?x # ?vts2 @ ?y
# ?vts3 @ [?x]))
unfolding simple-path-def by auto

have loop-free-p: loop-free p
using polygon-p p-is unfolding polygon-def simple-path-def by auto

have sublist-c1: sublist (?x # ?vts2 @ [?y]) vts
using  $\langle$ vts ! 0 # take (i - 1) (drop 1 vts) @ [vts ! i] = take (i + 1) (vts @ [vts
! 0]) $\rangle$  i-lt by auto
then have sublist-c1: sublist (?x # ?vts2 @ [?y]) (vts@[vts ! 0])
by (metis  $\langle$ vts ! 0 # take (i - 1) (drop 1 vts) @ [vts ! i] = take (i + 1) (vts
@ [vts ! 0]) $\rangle$  sublist-take)
then have loop-free ?c1
using sublist-is-loop-free p-is loop-free-p sublist-c1
by (metis One-nat-def Suc-1 Suc-eq-plus1 Suc-leI Suc-le-mono  $\langle$ vts ! 0 #
take (i - 1) (drop 1 vts) @ [vts ! i] = take (i + 1) (vts @ [vts ! 0]) $\rangle$  i-gt i-lt
length-append-singleton less-imp-le-nat take-i-is-loop-free)
then have simple-c1: simple-path ?c1
unfolding simple-path-def
using make-polygonal-path-gives-path by blast
have start-c1: pathstart ?c1 = ?x
using polygon-pathstart
by (metis Cons-eq-appendI list.discI nth-Cons-0 )
have finish-c1: pathfinish ?c1 = ?y
using polygon-pathfinish
by (metis Cons-eq-appendI diff-Suc-1 length-append-singleton list.discI nth-append-length)

have sublist-c2: sublist ([?y] @ ?vts3 @ [?x]) (vts@[vts ! 0])
by (metis  $\langle$ [vts ! i] @ drop i (drop 1 vts) @ [vts ! 0] = drop i (vts @ [vts ! 0]) $\rangle$ 
sublist-drop)
have i  $\leq$  length (tl vts) using i-lt by fastforce
then have loop-free ?c2
by (metis (no-types) Suc-1  $\langle$ [vts ! i] @ drop i (drop 1 vts) @ [vts ! 0] = drop
i (vts @ [vts ! 0]) $\rangle$   $\langle$ vts  $\neq$  [] $\rangle$  butlast-snoc drop-Suc drop-i-is-loop-free length-butlast
length-drop loop-free-p loop-free-reversepath p-is tl-append2)
then have simple-c2: simple-path ?c2
unfolding simple-path-def
using make-polygonal-path-gives-path

```

```

using path-imp-reversepath by blast
have start-c2: pathstart ?c2 = ?x
using polygon-pathfinish
by (metis (no-types, lifting) Nil-is-append-conv last-appendR last-conv-nth path-
start-reversepath polygon-pathfinish snoc-eq-iff-butlast)
have finish-c2: pathfinish ?c2 = ?y
using polygon-pathstart by auto

have path-image-int: path-image ?c1  $\subseteq$  path-image ?p
unfolding path-image-def
by (metis Un-upper1 p-is path-image-def path-image-p)
moreover have path-image ?p  $\cap$  path-image ?c3  $\subseteq$  {vts ! 0, vts ! i}
using have-cut unfolding is-polygon-cut-path-def
by (metis (no-types, lifting) Int-commute append-Cons append-is-Nil-conv cut-
path last-appendR last-conv-nth last-snoc not-Cons-self2 nth-Cons-0 polygon-pathfinish
polygon-pathstart set-eq-subset)
ultimately have vts-subset-c1c3: path-image ?c1  $\cap$  path-image ?c3  $\subseteq$  {?x, ?y}
by blast
have other-subset1: {vts ! 0, vts ! i}  $\subseteq$  path-image ?c1
using vertices-on-path-image by fastforce
have other-subset2: {vts ! 0, vts ! i}  $\subseteq$  path-image ?c3
unfolding assms using vertices-on-path-image by force
then have c1-inter-c3: path-image ?c1  $\cap$  path-image ?c3 = {vts ! 0, vts ! i}
using vts-subset-c1c3 other-subset1 other-subset2 by blast
then have path-image ?c1  $\cap$  path-image ?c3-rev = {pathstart ?c1, pathstart
?c3-rev}
by (metis rev-vts-path-image append-Cons append-Nil cutpath hd-conv-nth list.discI
list.sel(1) polygon-pathstart rev.simps(2) rev-rev-ident)

then have c1-inter-c3': path-image (make-polygonal-path (vts ! 0 # take (i -
1) (drop 1 vts) @ [vts ! i]))  $\cap$ 
path-image (make-polygonal-path (rev ([vts ! 0] @ cutvts @ [vts ! i])))
 $\subseteq$  {pathstart (make-polygonal-path (vts ! 0 # take (i - 1) (drop 1 vts) @ [vts !
i])),
pathstart (make-polygonal-path (rev ([vts ! 0] @ cutvts @ [vts ! i])))}
by blast
have last-is-head: last ?c3-rev-vts = hd ?c1-vts by auto
have vts-append: vts ! 0 # take (i - 1) (drop 1 vts) @ rev ([vts ! 0] @ cutvts @
[vts ! i]) =
(vts ! 0 # take (i - 1) (drop 1 vts) @ [vts ! i]) @
tl (rev ([vts ! 0] @ cutvts @ [vts ! i]))
by simp
have loop-free: loop-free (make-polygonal-path (vts ! 0 # take (i - 1) (drop 1
vts) @ [vts ! i]))  $\wedge$ 
loop-free (make-polygonal-path (rev ([vts ! 0] @ cutvts @ [vts ! i])))
by (metis Suc-eq-plus1 Suc-le-mono Zero-neq-Suc (vts ! 0 # take (i - 1) (drop
1 vts) @ [vts ! i]) = take (i + 1) (vts @ [vts ! 0])) cutpath diff-Suc-1 have-cut
i-gt i-lt is-polygon-cut-path-def length-append-singleton less-2-cases less-imp-le-nat

```

less-nat-zero-code linorder-le-less-linear loop-free-p p-is rev-vts-is-loop-free simple-path-def take-i-is-loop-free

have *last-is-head2*:
 $last (vts ! 0 \# take (i - 1) (drop 1 vts) @ [vts ! i]) =$
 $hd (rev ([vts ! 0] @ cutvts @ [vts ! i]))$ **by** *simp*
have *arcs*: $arc (make-polygonal-path (vts ! 0 \# take (i - 1) (drop 1 vts) @ [vts ! i])) \wedge$
 $arc (make-polygonal-path (rev ([vts ! 0] @ cutvts @ [vts ! i])))$
using *Nil-is-append-conv append-Cons constant-linepath-is-not-loop-free cutpath finish-c1 have-cut hd-conv-nth is-polygon-cut-path-def last-appendR last-conv-nth last-is-head last-is-head2 last-snoc list.sel(1) loop-free make-polygonal-path.simps(1) make-polygonal-path-gives-path polygon-pathfinish polygon-pathstart simple-path-def simple-path-imp-arc loop-free*
by (*smt (verit, ccfv-SIG)*)
then have *loop-free ?p1*
using *loop-free-append[of ?p1 ?p1-vts ?c1 ?c1-vts ?c3-rev ?c3-rev-vts,*
 $OF - - vts-append loop-free c1-inter-c3' - last-is-head2 arcs]$ **using**
last-is-head by blast

then have *simple-path ?p1*
unfolding *simple-path-def*
using *make-polygonal-path-gives-path by blast*
moreover have *closed-path ?p1*
using *polygon-pathstart polygon-pathfinish*
unfolding *closed-path-def*
using *elem-prop1 make-polygonal-path-gives-path*
by (*smt (verit, best) append-is-Nil-conv last-ConsR last-appendR last-conv-nth last-snoc list.discI nth-Cons-0 rev-append singleton-rev-conv*)
ultimately have *polygon-p1: polygon ?p1* **unfolding** *polygon-def polygonal-path-def*
by fastforce

have *path-image-int: path-image ?c2 \subseteq path-image (make-polygonal-path (vts @ [vts ! 0]))*
unfolding *path-image-def* **using** *path-image-p*
by (*simp add: p-is path-image-def*)
then have *vts-subset-c2c3: path-image ?c2 \cap path-image ?c3 \subseteq {?x, ?y}*
using *have-cut unfolding is-polygon-cut-path-def* **using** $\langle path-image (make-polygonal-path (vts @ [vts ! 0])) \cap path-image cutpath \subseteq \{vts ! 0, vts ! i\} \rangle$ **by auto**
have *other-subset3: {vts ! 0, vts ! i} \subseteq path-image ?c2*
using *vertices-on-path-image by fastforce*
have *other-subset4: {vts ! 0, vts ! i} \subseteq path-image ?c3*
unfolding *assms* **using** *vertices-on-path-image by fastforce*
have *c2-inter-c3: path-image ?c2 \cap path-image ?c3 = {vts ! 0, vts ! i}*
using *vts-subset-c2c3 other-subset3 other-subset4 by blast*
have *path-p2: path ?p2*
using *make-polygonal-path-gives-path by blast*
have *pathfinish ?p2 = vts ! 0*

```

using polygon-pathfinish
by (metis Nil-is-append-conv last-appendR last-conv-nth last-snoc list.discI)
then have closed-p2: closed-path ?p2
unfolding closed-path-def using polygon-pathstart
using path-p2 by auto

have ([vts ! 0] @ cutvts @ [vts ! i] @ drop i (drop 1 vts) @ [vts ! 0]) =
  ([vts ! 0] @ cutvts @ [vts ! i] @ tl ([vts ! i] @ drop i (drop 1 vts) @ [vts ! 0]))
by force
moreover have loop-free cutpath ∧
  loop-free (make-polygonal-path ([vts ! i] @ drop i (drop 1 vts) @ [vts ! 0]))
by (metis ‹loop-free (reversepath (make-polygonal-path ([vts ! i] @ drop i
(drop 1 vts) @ [vts ! 0])))› cutpath loop-free loop-free-reversepath rev-rev-ident
rev-vts-is-loop-free reversepath-reversepath)
moreover have path-image cutpath ∩ path-image (make-polygonal-path ([vts ! i]
@ drop i (drop 1 vts) @ [vts ! 0]))
  ⊆ {pathstart cutpath,
    pathstart (make-polygonal-path ([vts ! i] @ drop i (drop 1 vts) @ [vts ! 0]))}
using c2-inter-c3 cutpath polygon-pathstart by auto
moreover have last ([vts ! i] @ drop i (drop 1 vts) @ [vts ! 0]) ≠ hd ([vts ! 0]
@ cutvts @ [vts ! i]) →
  path-image cutpath ∩ path-image (make-polygonal-path ([vts ! i] @ drop i (drop
1 vts) @ [vts ! 0]))
  ⊆ {pathstart (make-polygonal-path ([vts ! i] @ drop i (drop 1 vts) @ [vts ! 0]))}
by simp
moreover have last ([vts ! 0] @ cutvts @ [vts ! i]) = hd ([vts ! i] @ drop i (drop
1 vts) @ [vts ! 0])
by simp
moreover have arc cutpath ∧ arc (make-polygonal-path ([vts ! i] @ drop i (drop
1 vts) @ [vts ! 0]))
by (metis (no-types, lifting) arc-simple-path arcs calculation(2) finish-c1 fin-
ish-c2 have-cut is-polygon-cut-path-def make-polygonal-path-gives-path pathfinish-reversepath
pathstart-reversepath simple-path-def start-c1 start-c2)
ultimately have loop-free ?p2
using loop-free-append[of ?p2 ?p2-vts ?c3 ?c3-vts ?c2' ?c2'-vts,
  OF - - ] using cutpath by blast
then have polygon-p2: polygon ?p2
using path-p2 closed-p2 unfolding polygon-def simple-path-def polygonal-path-def

by blast

have simple-c3: simple-path ?c3
using have-cut unfolding is-polygon-cut-path-def by meson
have start-c3: pathstart ?c3 = ?x unfolding assms using polygon-pathstart by
simp
have finish-c3: pathfinish ?c3 = ?y unfolding assms using polygon-pathfinish
by simp

```

have *pathstart cutpath* = ?*x* **using** *assms polygon-pathstart* **by force**
moreover have *pathfinish cutpath* = ?*y* **using** *assms polygon-pathfinish* **by simp**
ultimately have *vts-neg*: *vts ! 0* ≠ *vts ! i*
using *have-cut unfolding is-polygon-cut-path-def* **by force**
have *c1-inter-c2*: *path-image ?c1* ∩ *path-image ?c2* = {*vts ! 0*, *vts ! i*}
proof –
obtain *i* **where** *i1*: (?*x* # ?*vts2* @ [?*y*] = *take i (vts @ [vts!0])*) **and**
i2: ([?*y*] @ ?*vts3* @ [?*x*] = *drop (i-1) (vts @ [vts!0])*)
by (*metis* ⟨*[vts ! i] @ drop i (drop 1 vts) @ [vts ! 0] = drop i (vts @ [vts ! 0])*⟩,
⟨*vts ! 0 # take (i - 1) (drop 1 vts) @ [vts ! i] = take (i + 1) (vts @ [vts ! 0])*⟩,
add.commute add-diff-cancel-left)
moreover have *1*: *i* ≥ 1 ∧ *i* < *length (vts @ [vts!0])*
by (*metis* (*no-types*, *lifting*) *bot-nat-0.extremum less-one Nil-is-append-conv append-Cons calculation diff-is-0-eq drop-Cons' linorder-not-less list.inject not-Cons-self2 same-append-eq take-all vts-is vts-neg*)
moreover have *2*: ?*p* = *make-polygonal-path (vts @ [vts!0])* ∧ *loop-free ?p*
unfolding *polygon-of-def* **using** *p-is polygon-p unfolding polygon-def simple-path-def* **by blast**
ultimately have *path-image ?c1* ∩ *path-image (make-polygonal-path ([?y] @ ?vts3 @ [?x]))* ⊆ {*pathstart ?c1*, *pathstart (make-polygonal-path ([?y] @ ?vts3 @ [?x]))*}
using *loop-free-split-int[of ?p vts @ [vts!0] ?x # ?vts2 @ [?y] i [?y] @ ?vts3 @ [?x] ?c1 make-polygonal-path ([?y] @ ?vts3 @ [?x]) length (vts @ [vts!0])*,
OF 2 i1 i2 - - 1)
by presburger
moreover have *path-image ?c2* = *path-image (make-polygonal-path ([?y] @ ?vts3 @ [?x]))* **using** *path-image-reversepath* **by fast**
moreover have *pathstart (make-polygonal-path ([?y] @ ?vts3 @ [?x]))* = ?*y*
using *polygon-pathstart* **by auto**
moreover have *pathstart ?c1* = ?*x* **using** *polygon-pathstart* **by auto**
ultimately show ?*thesis*
using *other-subset1 other-subset3 subset-antisym* **by force**
qed

have *non-empty-inter*: *path-image ?c3* ∩ *inside(path-image ?c1 ∪ path-image ?c2)* ≠ {}
using *have-cut path-image-p p-is*
unfolding *is-polygon-cut-path-def path-inside-def*
by fastforce

have *p1-minus*: ((*path-image ?p1*) – (*path-image ?c3*)) = *path-image ?c1* – {?*x*, ?*y*}
using *c1-inter-c3 path-image-p1* **by blast**
have *p2-minus*: ((*path-image ?p2*) – (*path-image ?c3*)) = *path-image ?c2* – {?*x*, ?*y*}
using *c2-inter-c3 path-image-p2* **by auto**

then have *path-im-intersect-minus*: ((*path-image ?p1*) – (*path-image ?c3*)) ∩ ((*path-image ?p2*) – (*path-image (linepath ?x ?y)*)) = {}

```

using c1-inter-c2 p1-minus p2-minus
by blast
have  $((\text{path-image } ?p1) - (\text{path-image } ?c3)) \cup ((\text{path-image } ?p2) - (\text{path-image } ?c3)) \cup \{?x, ?y\} = ((\text{path-image } ?p1) - (\text{path-image } ?c3) \cup \{?x, ?y\}) \cup ((\text{path-image } ?p2) - (\text{path-image } ?c3) \cup \{?x, ?y\})$ 
by auto
then have  $((\text{path-image } ?p1) - (\text{path-image } (?c3))) \cup ((\text{path-image } ?p2) - (\text{path-image } (?c3))) \cup \{?x, ?y\} = ((\text{path-image } ?c1) - \{?x, ?y\} \cup \{?x, ?y\}) \cup ((\text{path-image } ?c2) - \{?x, ?y\} \cup \{?x, ?y\})$ 
using p1-minus p2-minus by simp
then have  $((\text{path-image } ?p1) - (\text{path-image } (?c3))) \cup ((\text{path-image } ?p2) - (\text{path-image } (?c3))) \cup \{?x, ?y\} = \text{path-image } ?c1 \cup \text{path-image } ?c2$ 
using other-subset1 other-subset3 by auto
then have path-im-intersect-union: path-image ?p = ((path-image ?p1) - (path-image ?c3)) \cup ((path-image ?p2) - (path-image ?c3)) \cup \{?x, ?y\}
using path-image-p p-is by auto

have  $\text{inside}(\text{path-image } ?c1 \cup \text{path-image } ?c3) \cap \text{inside}(\text{path-image } ?c2 \cup \text{path-image } ?c3) = \{\}$ 
using split-inside-simple-closed-curve-real2[OF simple-c1 start-c1 finish-c1 simple-c2 start-c2 finish-c2 simple-c3 start-c3 finish-c3 vts-neq c1-inter-c2 c1-inter-c3 c2-inter-c3 non-empty-inter]
by fast
then have empty-inter: path-inside ?p1 \cap path-inside ?p2 = \{\}
using path-image-p1 path-image-p2 unfolding path-inside-def
by force
have  $\text{inside}(\text{path-image } ?c1 \cup \text{path-image } ?c3) \cup \text{inside}(\text{path-image } ?c2 \cup \text{path-image } ?c3) \cup (\text{path-image } ?c3 - \{vts ! 0, vts ! i\}) = \text{inside}(\text{path-image } ?c1 \cup \text{path-image } ?c2)$ 
using split-inside-simple-closed-curve-real2[OF simple-c1 start-c1 finish-c1 simple-c2 start-c2 finish-c2 simple-c3 start-c3 finish-c3 vts-neq c1-inter-c2 c1-inter-c3 c2-inter-c3 non-empty-inter]
by fast
then have inside: path-inside ?p1 \cup path-inside ?p2 \cup (path-image ?c3 - \{?x, ?y\}) = path-inside p
using path-image-p1 path-image-p1 path-image-p unfolding path-inside-def
by (smt (z3) Diff-cancel Int-Un-distrib2 c1-inter-c2 c1-inter-c3 finish-c1 inf-commute inf-sup-absorb nonempty-simple-path-endless path-image-p2 simple-c1 start-c1)
have first-part: 0 < length vts \wedge i < length vts \wedge 0 < i
using assms
by auto
have second-part-helper: is-polygon-cut-path (vts @ [vts ! 0]) cutpath \wedge polygon ?p \wedge

```

$poly\!-\!p1 \wedge$
 $poly\!-\!p2 \wedge$
 $path\!-\!inside \ ?p1 \cap path\!-\!inside \ ?p2 = \{\}$ \wedge
 $path\!-\!inside \ ?p1 \cup path\!-\!inside \ ?p2 \cup (path\!-\!image \ (?c3) - \{?x, ?y\}) =$
 $path\!-\!inside \ p$
 $\wedge ((path\!-\!image \ ?p1) - (path\!-\!image \ (?c3))) \cap ((path\!-\!image \ ?p2) - (path\!-\!image$
 $(?c3))) = \{\}$
 $\wedge path\!-\!image \ ?p = ((path\!-\!image \ ?p1) - (path\!-\!image \ (?c3))) \cup ((path\!-\!image$
 $?p2) - (path\!-\!image \ (?c3))) \cup \{?x, ?y\}$
using *poly\!-\!p \ p\!-\!is \ poly\!-\!p1 \ poly\!-\!p2 \ empty\!-\!inter \ inside \ have\!-\!cut \ path\!-\!im\!-\!intersect\!-\!minus*
path\!-\!im\!-\!intersect\!-\!union
proof–
have $\{\} = path\!-\!image \ cutpath \cup path\!-\!image \ (make\!-\!polygonal\!-\!path \ (vts \ ! \ 0 \ \# \ take$
 $(i - 1) \ (drop \ 1 \ vts) \ @ \ [vts \ ! \ i])) \cap path\!-\!image \ (reversepath \ (make\!-\!polygonal\!-\!path$
 $([vts \ ! \ i] \ @ \ drop \ i \ (drop \ 1 \ vts) \ @ \ [vts \ ! \ 0]))) - path\!-\!image \ cutpath$
using *c1\!-\!inter\!-\!c2 \ c2\!-\!inter\!-\!c3* **by** *fastforce*
then have $\{\} = (path\!-\!image \ cutpath \cup path\!-\!image \ (make\!-\!polygonal\!-\!path \ (vts$
 $\ ! \ 0 \ \# \ take \ (i - 1) \ (drop \ 1 \ vts) \ @ \ [vts \ ! \ i])) \cap (path\!-\!image \ cutpath \cup path\!-\!image$
 $(reversepath \ (make\!-\!polygonal\!-\!path \ ([vts \ ! \ i] \ @ \ drop \ i \ (drop \ 1 \ vts) \ @ \ [vts \ ! \ 0]))) -$
 $path\!-\!image \ cutpath$
by *blast*
then show *?thesis*
using *empty\!-\!inter \ have\!-\!cut \ inside \ poly\!-\!p1 \ poly\!-\!p2 \ Int\!-\!Diff \ image\!-\!prop*
p\!-\!is \ path\!-\!im\!-\!intersect\!-\!union \ path\!-\!image\!-\!p2 \ poly\!-\!p
by *auto*
qed
have *vts\!-\!relation:* $(let \ vts1 = take \ 0 \ vts; \ vts2 = take \ (i - 0 - 1) \ (drop \ (Suc \ 0)$
 $vts);$
 $vts3 = drop \ (i - 0) \ (drop \ (Suc \ 0) \ vts); \ x = vts \ ! \ 0; \ y = vts \ ! \ i;$
 $p = make\!-\!polygonal\!-\!path \ (vts \ @ \ [vts \ ! \ 0]); \ p1 = make\!-\!polygonal\!-\!path \ (x \ \#$
 $vts2 \ @ \ ?c3\!-\!rev\!-\!vts);$
 $p2 = make\!-\!polygonal\!-\!path \ (?c3\!-\!vts \ @ \ vts3 \ @ \ [x]) \ in$
 $vts1 = [] \ \wedge \ vts2 = ?vts2 \ \wedge \ vts3 = ?vts3 \ \wedge \ p = ?p \ \wedge \ p1 = ?p1 \ \wedge \ p2 =$
 $?p2)$
by *simp*
have *second\!-\!part:* $(let \ vts1 = take \ 0 \ vts; \ vts2 = take \ (i - 0 - 1) \ (drop \ (Suc \ 0)$
 $vts);$
 $vts3 = drop \ (i - 0) \ (drop \ (Suc \ 0) \ vts); \ x = vts \ ! \ 0; \ y = vts \ ! \ i;$
 $p = make\!-\!polygonal\!-\!path \ (vts \ @ \ [vts \ ! \ 0]); \ p1 = make\!-\!polygonal\!-\!path \ (x \ \#$
 $vts2 \ @ \ ?c3\!-\!rev\!-\!vts);$
 $p2 = make\!-\!polygonal\!-\!path \ (vts1 \ @ \ ?c3\!-\!vts \ @ \ vts3 \ @ \ [vts \ ! \ 0])$
 $in \ is\!-\!polygon\!-\!cut\!-\!path \ (vts \ @ \ [vts \ ! \ 0]) \ cutpath \ \wedge$
 $poly\!-\!p \ \wedge$
 $poly\!-\!p1 \ \wedge$
 $poly\!-\!p2 \ \wedge$
 $path\!-\!inside \ p1 \cap path\!-\!inside \ p2 = \{\} \ \wedge$
 $path\!-\!inside \ p1 \cup path\!-\!inside \ p2 \cup (path\!-\!image \ cutpath - \{x, y\}) = path\!-\!inside$
 p
 $\wedge ((path\!-\!image \ p1) - (path\!-\!image \ (cutpath))) \cap ((path\!-\!image \ p2) - (path\!-\!image$

```

(cutpath)) = {} ∧
  path-image p = ((path-image p1) - (path-image (cutpath))) ∪ ((path-image
p2) - (path-image (cutpath))) ∪ {x, y}
  using second-part-helper vts-relation p-is
  by (metis self-append-conv2)
  show ?thesis
  unfolding is-polygon-split-path-def[of vts 0 i cutvts]
  using first-part second-part
  by (smt (verit, ccfv-threshold) append-Cons append-Nil cutpath rev.simps(2)
rev-append rev-is-Nil-conv)
qed

```

lemma *polygon-cut-path-to-split-path*:

```

fixes p :: R-to-R2
assumes polygon p
  p = make-polygonal-path (vts @ [vts ! 0])
  is-polygon-cut-path (vts @ [vts!0]) cutpath
  vts1 ≡ (take i vts)
  vts2 ≡ (take (j - i - 1) (drop (Suc i) vts))
  vts3 ≡ drop (j - i) (drop (Suc i) vts)
  x ≡ vts ! i
  y ≡ vts ! j
  cutpath = make-polygonal-path ([x] @ cutvts @ [y])
  i < length vts ∧ j < length vts ∧ i < j
  p1 ≡ make-polygonal-path (x#(vts2@[y] @ (rev cutvts) @ [x])) and
  p2 ≡ make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @ [(vts1 @
[x] ! 0])
  shows is-polygon-split-path vts i j cutvts
proof -
  let ?poly-vts-rot = rotate-polygon-vertices (vts @ [vts ! 0]) i
  let ?vts-rot = butlast ?poly-vts-rot
  let ?p-rot = make-polygonal-path ?poly-vts-rot
  let ?i-rot = j - i
  have rot-poly: polygon ?p-rot using assms(1) assms(2) rotation-is-polygon by
blast
  have i-rot: ?i-rot > 0 ∧ ?i-rot < length ?poly-vts-rot - 1
  using assms(10) rotate-polygon-vertices-same-length by fastforce
  have vtsi: vts ! i = ?poly-vts-rot ! 0
  using rotated-polygon-vertices[of ?poly-vts-rot vts @ [vts!0] i i]
  by (metis (no-types, lifting) One-nat-def Suc-1 assms(10) diff-self-eq-0 hd-conv-nth
last-snoc length-append-singleton less-imp-le-nat linorder-not-le not-less-eq-eq nth-append
take-all-iff take-eq-Nil)
  have vtsj: vts ! j = ?poly-vts-rot ! ?i-rot
  using rotated-polygon-vertices[of ?poly-vts-rot vts @ [vts!0] i j]
  by (smt (verit, ccfv-SIG) One-nat-def Suc-1 assms(10) butlast-snoc hd-append2
hd-conv-nth last-snoc leD length-append-singleton less-Suc-eq-le less-imp-le-nat not-less-eq-eq
nth-butlast take-all-iff take-eq-Nil)
  have is-polygon-cut-path ?poly-vts-rot cutpath
proof -

```

```

have ?poly-vts-rot ! 0 ≠ ?poly-vts-rot ! ?i-rot
  using assms(3) unfolding is-polygon-cut-path-def using vtsi vtsj
  using append-Cons append-is-Nil-conv assms(7) assms(8) assms(9) last-appendR
last-conv-nth polygon-pathfinish polygon-pathstart
  by force
  moreover have {?poly-vts-rot ! 0, ?poly-vts-rot ! ?i-rot} ⊆ set (?poly-vts-rot
@ [?poly-vts-rot ! 0])
    using assms(3) unfolding is-polygon-cut-path-def using i-rot vtsi vtsj by
fastforce
    moreover have path-image cutpath ∩ path-image ?p-rot = {?poly-vts-rot ! 0,
?poly-vts-rot ! ?i-rot}
      using polygon-vts-rotation vtsi vtsj assms(3) is-polygon-cut-path-def
      by (metis (no-types, lifting) append.assoc append-Cons assms(7) assms(8)
assms(9) last-conv-nth nth-Cons-0 polygon-pathfinish polygon-pathstart snoc-eq-iff-butlast)
      moreover have path-image cutpath ∩ path-inside (?p-rot) ≠ {}
        using vtsi vtsj assms(3) polygon-vts-rotation
        unfolding is-polygon-cut-path-def path-inside-def by metis
        ultimately show ?thesis
        unfolding is-polygon-cut-path-def
        using rot-poly assms(3) is-polygon-cut-path-def rotate-polygon-vertices-same-set
vtsi vtsj
        by (metis polygon-vts-rotation)
    qed
  then have rot-cut: is-polygon-cut-path (?vts-rot @ [?vts-rot!0]) cutpath
    by (metis butlast-snoc rotate-polygon-vertices-def)
    have rot-cut-butlast: make-polygonal-path ?poly-vts-rot = make-polygonal-path
(?vts-rot @ [?vts-rot!0])
      by (metis butlast-snoc rotate-polygon-vertices-def)
      have split-rot: is-polygon-split-path ?vts-rot 0 ?i-rot cutvts
        using rot-cut rot-cut-butlast
        by (smt (verit, ccfv-SIG) assms(7) assms(8) assms(9) dual-order.strict-trans
i-rot is-polygon-cut-path-def length-butlast nth-butlast polygon-cut-path-to-split-path-vtx0
vtsi vtsj)

let ?vts1-rot = take 0 ?vts-rot
let ?vts2-rot = take (j - i - 0 - 1) (drop (Suc 0) ?vts-rot)
let ?vts3-rot = drop (j - i - 0) (drop (Suc 0) ?vts-rot)
let ?x-rot = ?vts-rot ! 0
let ?y-rot = ?vts-rot ! (j - i)
let ?p1-rot-vts = ?x-rot # ?vts2-rot @ [?y-rot] @ (rev cutvts) @ [?x-rot]
let ?p1-rot = make-polygonal-path ?p1-rot-vts
let ?p2-rot-vts = ?vts1-rot @ [?x-rot] @ cutvts @ [?y-rot] @ ?vts3-rot @ [?vts-rot
! 0]
let ?p2-rot = make-polygonal-path ?p2-rot-vts

let ?p1-vts = x # vts2 @ [y] @ (rev cutvts) @ [x]
let ?p2-vts = vts1 @ [x] @ cutvts @ [y] @ vts3 @ [(vts1 @ [x]) ! 0]

have p2-firstlast: hd ?p2-vts = last ?p2-vts

```

by (*metis* (*no-types*, *lifting*) *append-is-Nil-conv* *append-self-conv2* *hd-append2* *hd-conv-nth* *last-appendR* *last-snoc* *list.discI* *list.sel(1)*)

have $\text{length } (\text{drop } (\text{Suc } i) \text{ vts}) = \text{length } \text{vts} - i - 1$
by *simp*

then have $\text{len-prop: } \text{length } (\text{drop } (\text{Suc } i) \text{ vts}) \geq j - i - 1$
using *assms(9)* *assms(10)* *diff-le-mono* *less-or-eq-imp-le* **by** *presburger*

have $\text{drop-take: } \text{rotate } i \text{ vts} = \text{drop } i \text{ vts} @ \text{take } i \text{ vts}$
using *rotate-drop-take[of i vts]* *assms(10)* *mod-less* **by** *presburger*

then have $\text{drop-take-suc: } \text{drop } (\text{Suc } 0) (\text{rotate } i \text{ vts}) = \text{drop } (\text{Suc } i) \text{ vts} @ \text{take } i \text{ vts}$
using *assms(10)* **by** *simp*

then have $\text{take } (j - \text{Suc } i) (\text{drop } (\text{Suc } 0) (\text{rotate } i \text{ vts})) = \text{take } (j - \text{Suc } i) (\text{drop } (\text{Suc } i) \text{ vts})$
using *len-prop* **by** *force*

then have $\text{vts2: } \text{take } (j - i - 0 - 1) (\text{drop } (\text{Suc } 0) (\text{butlast } (\text{rotate-polygon-vertices } (\text{vts} @ [\text{vts} ! 0]) i))) = \text{vts2}$
using *assms(5)* **unfolding** *rotate-polygon-vertices-def*
by (*metis* *Suc-eq-plus1* *butlast-snoc* *diff-diff-left* *diff-zero*)

have $xy: ?x\text{-rot} = x \wedge ?y\text{-rot} = y$
using *vtsi* *vtsj* *assms* **by** (*metis* *is-polygon-split-path-def* *nth-butlast* *split-rot*)

moreover have $\text{path-image } p = \text{path-image } ?p\text{-rot}$
using *assms(1)* *assms(2)* *polygon-vts-arb-rotation* **by** *auto*

moreover then have $\text{path-inside } p = \text{path-inside } ?p\text{-rot}$ **unfolding** *path-inside-def*
by *simp*

moreover have $?p1\text{-rot-vts} = ?p1\text{-vts}$ **using** *xy* *vts2* **by** *presburger*

moreover then have $\text{path-image } p1 = \text{path-image } ?p1\text{-rot}$ **using** *assms* **by** *argo*

moreover then have $\text{path-inside } p1 = \text{path-inside } ?p1\text{-rot}$ **unfolding** *path-inside-def*
by *argo*

moreover have $\text{polygon } p1$
using *calculation* *split-rot* *assms(11)* **unfolding** *is-polygon-split-path-def*
by (*smt* (*verit*, *ccfv-SIG*) *vts2*)

moreover have $?p2\text{-rot-vts} = \text{rotate-polygon-vertices } ?p2\text{-vts } i$
proof–

have $\text{butlast } (\text{vts1} @ [x] @ \text{cutvts} @ [y] @ \text{vts3} @ [(\text{vts1} @ [x]) ! 0])$
 $= \text{vts1} @ [x] @ \text{cutvts} @ [y] @ \text{vts3}$
by (*simp* *add: butlast-append*)

also have $\text{rotate } i \dots = [x] @ \text{cutvts} @ [y] @ \text{vts3} @ \text{vts1}$
using *assms(4)*

by (*metis* (*no-types*, *lifting*) *drop-take* *add-diff-cancel-right'* *append.assoc* *assms(10)* *diff-diff-cancel* *length-append* *length-drop* *length-rotate* *less-imp-le-nat* *rotate-append*)

finally have $\text{rotate-polygon-vertices } ?p2\text{-vts } i = [x] @ \text{cutvts} @ [y] @ \text{vts3} @ \text{vts1} @ [x]$
unfolding *rotate-polygon-vertices-def* **by** *simp*

```

moreover have ?vts3-rot = vts3 @ vts1
  using assms(4,6) unfolding rotate-polygon-vertices-def
    by (smt (verit, del-insts) One-nat-def Suc-diff-Suc Suc-leI drop-take-suc
assms(10) butlast-snoc diff-is-0-eq diff-zero drop0 drop-append i-rot le-add-diff-inverse
len-prop length-drop nat-less-le)
  ultimately show ?thesis by (simp add: xy)
qed
moreover then have polygon p2
  using unrotation-is-polygon[of ?p2-vts i p2] split-rot assms(12) p2-firstlast
  unfolding is-polygon-split-path-def
  by (smt (verit) append.assoc)
moreover then have path-image p2 = path-image (?p2-rot)
  using assms(12) polygon-vts-arb-rotation calculation by auto
moreover then have path-inside p2 = path-inside ?p2-rot unfolding path-inside-def
by presburger

```

```

ultimately show is-polygon-split-path vts i j cutvts
  using split-rot unfolding is-polygon-split-path-def
  using One-nat-def assms bot-nat-0.not-eq-extremum butlast-snoc hd-append2
hd-conv-nth hd-take le-add2 length-0-conv length-Cons length-append length-butlast
nth-append-length rot-cut-butlast rotate-polygon-vertices-same-length take-eq-Nil
  by (smt (verit) append.assoc butlast-conv-take have-wraparound-vertex is-polygon-cut-path-def
rotate-polygon-vertices-same-set)
qed

```

lemma *good-polygonal-path-implies-polygon-split-path:*

```

assumes polygon p
assumes p = make-polygonal-path (vts @ [vts!0])
assumes good-polygonal-path v1 cutvts v2 (vts @ [vts!0])
assumes i < length vts ∧ j < length vts
assumes vts ! i = v1
assumes vts ! j = v2
assumes i < j
shows is-polygon-split-path vts i j cutvts
proof –
let ?cutpath = make-polygonal-path ([v1] @ cutvts @ [v2])
let ?p-path = make-polygonal-path (vts @ [vts!0])
have linepath-subset: path-image ?cutpath ⊆ path-inside ?p-path ∪ {v1, v2}
  using assms(3) unfolding good-polygonal-path-def by meson
have linepath-ends: pathstart ?cutpath = v1 ∧ pathfinish ?cutpath = v2
  using polygon-pathfinish polygon-pathstart by force
then have vs-subset1: {v1, v2} ⊆ path-image ?cutpath
  using vertices-on-path-image by fastforce
have vs-subset2: {v1, v2} ⊆ path-image (make-polygonal-path (vts @ [vts ! 0]))
  using assms(4-6) vertices-on-path-image[of vts]
  using vertices-on-path-image by fastforce
have path-inside ?p-path ∩ path-image ?p-path = {}
using inside-outside-polygon[OF assms(1)] assms(2) unfolding inside-outside-def
by blast

```

```

then have linepath-path: path-image ?cutpath  $\cap$  path-image (make-polygonal-path
(vts @ [vts ! 0])) = {v1, v2}
  using linepath-subset vs-subset1 vs-subset2
  by blast
have ?cutpath (5 / 10)  $\in$  path-image ?cutpath
  unfolding path-image-def by auto
have v1-neq-v2: v1  $\neq$  v2
  using assms(3) unfolding good-polygonal-path-def
  by fastforce
have not-v1: ?cutpath (0.5::real) = v1  $\implies$  False
proof –
  assume *: ?cutpath (0.5::real) = v1
  then have ?cutpath (0.5::real) = ?cutpath 0
    using linepath-ends unfolding pathstart-def by simp
  moreover have loop-free ?cutpath using assms unfolding good-polygonal-path-def
by metis
  ultimately show False unfolding loop-free-def by fastforce
qed
have not-v2: ?cutpath (0.5::real) = v2  $\implies$  False
proof –
  assume *: ?cutpath (0.5::real) = v2
  then have ?cutpath (0.5::real) = ?cutpath 1
    using linepath-ends unfolding pathfinish-def by simp
  moreover have loop-free ?cutpath using assms unfolding good-polygonal-path-def
by metis
  ultimately show False unfolding loop-free-def by fastforce
qed
then have ?cutpath (0.5::real)  $\neq$  v1  $\wedge$  ?cutpath (0.5::real)  $\neq$  v2
  using not-v1 not-v2 by auto
then have linepath-inside: path-image ?cutpath  $\cap$  path-inside (make-polygonal-path
(vts @ [vts ! 0]))  $\neq$  {}
  using linepath-subset
  using  $\langle$ ?cutpath (5 / 10)  $\in$  path-image ?cutpath $\rangle$  by blast
have is-polygon-cut-path (vts @ [vts!0]) ?cutpath
  using assms(3) assms(1–2) unfolding good-polygonal-path-def is-polygon-cut-path-def
  using linepath-path linepath-inside
  by (metis linepath-ends make-polygonal-path-gives-path simple-path-def)
then show ?thesis using polygon-cut-path-to-split-path assms by blast
qed

```

lemma *good-path-iff*:

```

good-linepath a b vts  $\longleftrightarrow$  good-polygonal-path a [] b vts
unfolding good-linepath-def good-polygonal-path-def
using linepath-loop-free by auto

```

lemma *polygon-cut-iff*: *is-polygon-cut* (*vts* @ [*vts*!0]) (*vts*!*i*) (*vts*!*j*)
 \longleftrightarrow *is-polygon-cut-path* (*vts* @ [*vts*!0]) (*linepath* (*vts*!*i*) (*vts*!*j*))

```

unfolding is-polygon-cut-def is-polygon-cut-path-def
by (metis pathfinish-linepath pathstart-linepath simple-path-linepath)

lemma polygon-split-iff: is-polygon-split vts i j  $\longleftrightarrow$  is-polygon-split-path vts i j []
unfolding is-polygon-split-def is-polygon-split-path-def
by (smt (verit, cfv-threshold) append-Cons append-Nil make-polygonal-path.simps(3)
polygon-cut-iff rev.simps(1))

lemma polygon-cut-to-split-vtx0:
fixes p :: R-to-R2
assumes polygon-p: polygon p and
i-gt: i > 0 and
i-lt: i < length vts and
p-is: p = make-polygonal-path (vts @ [vts ! 0]) and
have-cut: is-polygon-cut (vts @ [vts!0]) (vts!0) (vts!i)
shows is-polygon-split vts 0 i
using have-cut i-gt i-lt p-is polygon-cut-path-to-split-path-vtx0 polygon-cut-iff poly-
gon-p polygon-split-iff
by force

lemma polygon-cut-to-split:
fixes p :: R-to-R2
assumes is-polygon-cut (vts @ [vts!0]) (vts!i) (vts!j)
i < length vts  $\wedge$  j < length vts  $\wedge$  i < j
shows is-polygon-split vts i j
by (metis append-Cons append-Nil assms is-polygon-cut-def make-polygonal-path.simps(3)
polygon-cut-path-to-split-path polygon-cut-iff polygon-split-iff)

lemma good-linepath-implies-polygon-split:
assumes polygon p
assumes p = make-polygonal-path (vts @ [vts!0])
assumes good-linepath v1 v2 (vts @ [vts!0])
assumes i < length vts  $\wedge$  j < length vts
assumes vts ! i = v1
assumes vts ! j = v2
assumes i < j
shows is-polygon-split vts i j
using assms good-path-iff good-polygonal-path-implies-polygon-split-path polygon-split-iff
by auto

end
theory Triangle-Lemmas
imports
Polygon-Convex-Lemmas
Integral-Matrix
Affine-Arithmetic.Floatarith-Expression
HOL-Analysis.Topology-Euclidean-Space
HOL-Analysis.Equivalence-Lebesgue-Henstock-Integration
HOL-Analysis.Inner-Product

```

HOL-Analysis.Line-Segment
HOL-Analysis.Convex-Euclidean-Space
HOL-Analysis.Change-Of-Vars

begin

20 Triangles

definition *elem-triangle* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{bool}$ **where**

elem-triangle $a\ b\ c \iff$
 $\neg \text{collinear } \{a, b, c\}$
 $\wedge \text{integral-vec } a \wedge \text{integral-vec } b \wedge \text{integral-vec } c$
 $\wedge \{x. x \in \text{convex hull } \{a, b, c\} \wedge \text{integral-vec } x\} = \{a, b, c\}$

definition *triangle-mat* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2$ **where**

triangle-mat $a\ b\ c = \text{transpose } (\text{vector } [b - a, c - a])$

definition *triangle-linear* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow (\text{real}^2 \Rightarrow \text{real}^2)$
where

triangle-linear $a\ b\ c = (\lambda x. (\text{triangle-mat } a\ b\ c) *v x)$

definition *triangle-affine* :: $\text{real}^2 \Rightarrow \text{real}^2 \Rightarrow \text{real}^2 \Rightarrow (\text{real}^2 \Rightarrow \text{real}^2)$ **where**

triangle-affine $a\ b\ c = (\lambda x. a + (\text{triangle-mat } a\ b\ c) *v x)$

abbreviation *unit-square* \equiv

$(\text{convex hull } \{\text{vector } [0, 0], \text{vector } [0, 1], \text{vector } [1, 1], \text{vector } [1, 0]\}) :: ((\text{real}^2) \text{ set})$

abbreviation *unit-triangle* \equiv

$(\text{convex hull } \{\text{vector } [0, 0], \text{vector } [1, 0], \text{vector } [0, 1]\}) :: ((\text{real}^2) \text{ set})$

abbreviation *unit-triangle'* \equiv

$(\text{convex hull } \{\text{vector } [1, 1], \text{vector } [1, 0], \text{vector } [0, 1]\}) :: ((\text{real}^2) \text{ set})$

lemma *triangle-inside-is-convex-hull-interior*:

assumes *polygon-of* $p\ [a, b, c, a]$

shows *path-inside* $p = \text{interior } (\text{convex hull } \{a, b, c\})$

proof –

have *path-image* $p = \text{closed-segment } a\ b \cup \text{closed-segment } b\ c \cup \text{closed-segment } c\ a$

proof –

have *path-image* $(\text{linepath } a\ b) = \text{closed-segment } a\ b$ **by** *simp*

moreover have *path-image* $(\text{linepath } b\ c) = \text{closed-segment } b\ c$ **by** *simp*

moreover have *path-image* $(\text{linepath } c\ a) = \text{closed-segment } c\ a$ **by** *simp*

moreover have *path-image* $p = \text{path-image } (\text{linepath } a\ b) \cup \text{path-image } (\text{linepath } b\ c) \cup \text{path-image } (\text{linepath } c\ a)$

using *calculation* *assms*(1) **unfolding** *polygon-of-def* *make-polygonal-path.simps*

by (*simp* *add: path-image-join sup-assoc*)

ultimately show *?thesis* **by** *simp*

qed
moreover have $DIM((real, 2) vec) = 2$ **by** *simp*
ultimately show *?thesis using inside-of-triangle[of a b c] unfolding path-inside-def*
by *presburger*
qed

lemma *triangle-is-convex*:
assumes $p = make_triangle\ a\ b\ c$ **and** $\neg collinear\ \{a, b, c\}$
shows *convex (path-inside p) (is convex ?s)*
using *triangle-inside-is-convex-hull-interior assms(1) assms(2)*
using *make-triangle-def polygon-of-def triangle-is-polygon*
by *auto*

lemma *affine-comp-linear-trans*: $triangle_affine\ a\ b\ c = (\lambda x. x + a) \circ (triangle_linear\ a\ b\ c)$
apply (*simp add: triangle-affine-def triangle-linear-def*)
by *auto*

lemma *triangle-linear-der*:
fixes $a\ b\ c :: real^2$
defines $T \equiv triangle_linear\ a\ b\ c$
shows (*T has-derivative T*) (at x)
proof –
have *linear T using T-def by (simp add: triangle-linear-def)*
then have *bounded-linear T by (simp add: linear-linear)*
thus *?thesis using bounded-linear-imp-has-derivative by blast*
qed

lemma *triangle-affine-der*:
fixes $a\ b\ c :: real^2$
assumes $S \in sets\ lebesgue$ **and** $x \in S$
defines $A \equiv triangle_affine\ a\ b\ c$ **and** $T \equiv triangle_linear\ a\ b\ c$
shows $x \in S \implies (A\ has_derivative\ T)$ (at x within S)
proof –
assume $xin: x \in S$
let $?trans = \lambda x :: real^2. x + a$
have $comp: (?trans \circ T) = (\lambda x. (T\ x) + a)$
by *auto*
have $\forall x. A\ x = (?trans \circ T)\ x$ **unfolding** *A-def T-def using affine-comp-linear-trans*
by *auto*
moreover then have $Ax-is: (\bigwedge x. x \in S \implies A\ x = ((\lambda x. x + a) \circ T)\ x)$
by *auto*
moreover have *trans-der: (?trans has-derivative id) (at x within S)*
by (*metis (full-types) add commute assms(2) eq-id-iff has-derivative-transform shift-has-derivative-id*)
moreover have *Tder: (T has-derivative T) (at x within S) using triangle-linear-der*
by (*simp add: T-def bounded-linear-imp-has-derivative triangle-linear-def*)
moreover have *comp-der: ((?trans \circ T) has-derivative T) (at x within S)*
using *has-derivative-add-const[OF Tder] comp*

```

    by simp
  ultimately show (A has-derivative T) (at x within S)
    using triangle-affine-def triangle-linear-def affine-comp-linear-trans o-apply
  add.commute vector-derivative-chain-within assms(2) has-derivative-add-const has-derivative-transform
  A-def T-def
    by force
  qed

```

lemma *triangle-linear-inj*:

```

  fixes a b c :: real^2
  assumes ¬ collinear {a, b, c}
  defines L ≡ triangle-linear a b c
  shows inj L
proof -
  let ?M = triangle-mat a b c
  let ?m-11 = (b - a)$1
  let ?m-12 = (c - a)$1
  let ?m-21 = (b - a)$2
  let ?m-22 = (c - a)$2
  have det ?M = ?m-11*?m-22 - ?m-12*?m-21
    unfolding triangle-mat-def
    by (metis det-2 det-transpose mult.commute vector-2(1) vector-2(2))
  moreover have ?m-11*?m-22 ≠ ?m-12*?m-21
  proof (rule ccontr)
    assume ¬ ?m-11*?m-22 ≠ ?m-12*?m-21
    then have eq: ?m-11*?m-22 = ?m-12*?m-21 by simp
    { assume *: ?m-21 = 0 ∧ ?m-22 ≠ 0
      then have ?m-11 = 0 using eq by simp
      then have ?m-11 = 0 ∧ ?m-21 = 0 using * by auto
      then have b - a = 0 by (metis (no-types, opaque-lifting) exhaust-2 vec-eq-iff
zero-index)
      then have collinear {a, b, c} by simp
      then have False using assms by fastforce
    } moreover
    { assume *: ?m-21 ≠ 0 ∧ ?m-22 = 0
      then have ?m-12 = 0 using eq by simp
      then have ?m-12 = 0 ∧ ?m-22 = 0 using * by auto
      then have c - a = 0 by (metis (no-types, opaque-lifting) exhaust-2 vec-eq-iff
zero-index)
      then have collinear {a, b, c} by (simp add: collinear-3-eq-affine-dependent)
      then have False using assms by fastforce
    } moreover
    { assume *: ?m-21 = 0 ∧ ?m-22 = 0
      { assume ?m-11 = 0
        then have b - a = 0 using *
          by (metis (no-types, opaque-lifting) exhaust-2 vec-eq-iff zero-index)
        then have False using assms(1) by auto
      } moreover
      { assume ?m-11 ≠ 0

```

then obtain k where $?m-12 = k * ?m-11$ using *nonzero-divide-eq-eq* by
blast
moreover have $?m-22 = k * ?m-21$ using $*$ by *auto*
ultimately have $c - a = k *_R (b - a)$
by (*smt (verit, del-insts) exhaust-2 real-scaleR-def vec-eq-iff vector-scaleR-component*)
then have *collinear* $\{a, b, c\}$
using *vec-diff-scale-collinear*[of c a k b] by (*simp add: insert-commute*)
then have *False* using *assms(1)* by *fastforce*
}
ultimately have *False* using *assms* by *fastforce*
} moreover
{ assume $*$: $?m-21 \neq 0 \wedge ?m-22 \neq 0$
then have $?m-11 / ?m-21 = ?m-12 / ?m-22$ using *eq frac-eq-eq* by *blast*
then obtain m where $?m-11 = m * ?m-12 \wedge ?m-21 = m * ?m-22$
using *nonzero-divide-eq-eq* $*$
by (*metis (no-types, lifting) mult.commute times-divide-eq-left*)
then have $b - a = m * s (c - a)$
by (*smt (verit, del-insts) exhaust-2 vec-eq-iff vector-smult-component*)
then have $b - a = m *_R (c - a)$ by (*simp add: scalar-mult-eq-scaleR*)
then have *collinear* $\{a, b, c\}$ using *vec-diff-scale-collinear* by *auto*
then have *False* using *assms* by *auto*
}
ultimately show *False* by *fastforce*
qed
ultimately have $\det ?M \neq 0$ by *linarith*
thus *?thesis* by (*simp add: L-def inj-matrix-vector-mult invertible-det-nz triangle-linear-def*)
qed

lemma *triangle-affine-inj*:
fixes $a b c :: \text{real}^2$
assumes $\neg \text{collinear } \{a, b, c\}$
defines $A \equiv \text{triangle-affine } a b c$
shows *inj* A
proof –
have *inj* (*triangle-linear* $a b c$) using *triangle-linear-inj*[of $a b c$] *assms* by *auto*
moreover have *inj* $(\lambda x. x + a)$ by *simp*
moreover have $A = (\lambda x. x + a) \circ (\text{triangle-linear } a b c)$
by (*simp add: A-def affine-comp-linear-trans*)
ultimately show *?thesis* using *inj-compose* by *blast*
qed

lemma *triangle-linear-integrable*:
fixes $a b c :: \text{real}^2$
assumes $S \in \text{lmeasurable}$
defines $T \equiv \text{triangle-linear } a b c$
shows $(\lambda x. \text{abs } (\det (\text{matrix } (T)))) \text{integrable-on } S$ (is $(\lambda x. ?c) \text{integrable-on } S$)
using *integrable-on-const*[of S $?c$] *assms(1)* by *blast*

lemma *measure-differentiable-image-eq-affine*:
fixes $a\ b\ c :: \text{real}^2$
defines $A \equiv \text{triangle-affine } a\ b\ c$ **and** $T \equiv \text{triangle-linear } a\ b\ c$
assumes $S \in \text{lmeasurable}$ **and** $\neg \text{collinear } \{a, b, c\}$
shows $\text{measure lebesgue } (A \text{ ` } S) = \text{integral } S (\lambda x. \text{abs } (\text{det } (\text{matrix } T)))$
proof –
have $\bigwedge x. x \in S \implies (A \text{ has-derivative } T) \text{ (at } x \text{ within } S)$
using *triangle-affine-der A-def T-def assms(3)* **by** *blast*
moreover **have** *inj-on A S*
using *A-def assms(3) assms(4) triangle-affine-inj inj-on-subset* **by** *blast*
moreover **have** $(\lambda x. \text{abs } (\text{det } (\text{matrix } (T)))) \text{ integrable-on } S$
by *(simp add: T-def assms(3) triangle-linear-integrable)*
ultimately show *?thesis*
using *measure-differentiable-image-eq[of - - \lambda x. T] assms(3)* **by** *blast*
qed

lemma *triangle-affine-img*:
fixes $a\ b\ c :: \text{real}^2$
defines $A \equiv \text{triangle-affine } a\ b\ c$
shows $\text{convex hull } \{a, b, c\} = A \text{ ` unit-triangle}$
proof –
let $?O = (\text{vector } [0, 0]) :: \text{real}^2$
let $?e1 = (\text{vector } [1, 0]) :: \text{real}^2$
let $?e2 = (\text{vector } [0, 1]) :: \text{real}^2$

let $?translate-a = \lambda x. x + a$

let $?T = \text{triangle-linear } a\ b\ c$

define al **where** $al = ?T ?O$
define bl **where** $bl = ?T ?e1$
define cl **where** $cl = ?T ?e2$

have $a: a = ?translate-a\ al$
proof –
have $al = ?O$
by *(simp add: al-def mat-vec-mult-2 triangle-linear-def)*
then show *?thesis*
by *(metis (no-types, opaque-lifting) add-0 mat-vec-mult-2 matrix-vector-mult-0 mult-zero-right zero-index)*
qed
have $b: b = ?translate-a\ bl$
proof –
have $col1: \text{column } 1 \text{ (triangle-mat } a\ b\ c) = b - a$
by *(metis column-transpose row-def triangle-mat-def vec-lambda-eta vector-2(1))*
then have $bl = b - a$
using *bl-def unfolding triangle-linear-def triangle-mat-def matrix-vector-mult-def*
using *matrix-vector-mult-basis[of triangle-mat a b c 1]*

```

    by (simp add: col1 axis-def bl-def mat-vec-mult-2 triangle-linear-def)
  then show ?thesis by simp
qed
have c: c = ?translate-a cl
proof-
  have col2: column 2 (triangle-mat a b c) = c - a
  by (metis column-transpose row-def triangle-mat-def vec-lambda-eta vector-2(2))
  then have cl = c - a
  using cl-def unfolding triangle-linear-def triangle-mat-def matrix-vector-mult-def
  using matrix-vector-mult-basis[of triangle-mat a b c 2]
  by (simp add: col2 axis-def cl-def mat-vec-mult-2 triangle-linear-def)
  then show ?thesis by simp
qed

have linear ?T using triangle-linear-def by force
then have ?T ' unit-triangle = convex hull {al, bl, cl}
  using convex-hull-linear-image al-def bl-def cl-def by force
also have ?translate-a ' ... = convex hull {a, b, c}
  using a b c convex-hull-translation[of a {al, bl, cl}]
  by (metis (no-types, lifting) add commute image-cong image-empty image-insert)
finally have ?translate-a ' (?T ' unit-triangle) = convex hull {a, b, c} .
moreover have ?translate-a o ?T = A unfolding A-def using affine-comp-linear-trans
by auto
ultimately show ?thesis by fastforce
qed

lemma triangle-affine-e1-e2:
  fixes a b c :: real^2
  defines A ≡ triangle-affine a b c
  shows (triangle-affine a b c) (vector [0, 0]) = a
        (triangle-affine a b c) (vector [1, 0]) = b
        (triangle-affine a b c) (vector [0, 1]) = c
proof-
  let ?M = triangle-mat a b c
  let ?L = triangle-linear a b c
  let ?A = triangle-affine a b c
  let ?O = (vector [0, 0])::(real^2)
  let ?e1 = (vector [1, 0])::(real^2)
  let ?e2 = (vector [0, 1])::(real^2)

  show ?A ?O = a
  unfolding triangle-affine-def triangle-mat-def
  by (metis (no-types, opaque-lifting) add.right-neutral diff-self mult-zero-right
scaleR-left-diff-distrib transpose-matrix-vector vec-scaleR-2 vector-matrix-mult-0)
  show ?A ?e1 = b
proof-
  have ?L ?e1 = ?M *v ?e1 unfolding triangle-linear-def by blast
  also have ... = vector [1*(?M$1$1) + 0*(?M$1$2), 1*(?M$2$1) + 0*(?M$2$2)]

```

```

    unfolding triangle-linear-def triangle-mat-def
    using mat-vec-mult-2 by force
  also have ... = vector [1*(b - a)$1 + 0*(?M$1$2), 1*(b - a)$2 + 0*(?M$2$2)]
    unfolding triangle-mat-def transpose-def by simp
  also have ... = vector [(b - a)$1, (b - a)$2] by argo
  also have ... = b - a
    by (smt (verit) exhaust-2 vec-eq-iff vector-2(1) vector-2(2))
  finally show ?thesis unfolding triangle-affine-def triangle-linear-def by simp
qed
show ?A ?e2 = c
proof-
  have ?L ?e2 = ?M *v ?e2 unfolding triangle-linear-def by blast
  also have ... = vector [0*(?M$1$1) + 1*(?M$1$2), 0*(?M$2$1) + 1*(?M$2$2)]
    unfolding triangle-linear-def triangle-mat-def
    using mat-vec-mult-2 by force
  also have ... = vector [0*(?M$1$1) + 1*(c - a)$1, 0*(?M$2$1) + 1*(c -
a)$2]
    unfolding triangle-mat-def transpose-def by simp
  also have ... = vector [(c - a)$1, (c - a)$2] by argo
  also have ... = c - a
    by (smt (verit) exhaust-2 vec-eq-iff vector-2(1) vector-2(2))
  finally show ?thesis unfolding triangle-affine-def triangle-linear-def by simp
qed
qed

lemma triangle-measure-integral-of-det:
  fixes a b c :: real^2
  defines S ≡ convex hull {a, b, c}
  assumes ¬ collinear {a, b, c}
  shows measure lebesgue S =
    integral unit-triangle (λ(x::real^2). abs (det (matrix (triangle-linear a b
c))))
proof-
  let ?A = triangle-affine a b c
  let ?T = triangle-linear a b c

  have bounded unit-triangle by (simp add: finite-imp-bounded-convex-hull)
  then have lmeasurable-S: unit-triangle ∈ lmeasurable
    using bounded-set-imp-lmeasurable measurable-convex by blast

  have S = ?A ‘ unit-triangle using S-def triangle-affine-img by blast
  then have measure lebesgue S = measure lebesgue (?A ‘ unit-triangle) by blast
  moreover have
    measure lebesgue (?A ‘ unit-triangle)
    = integral unit-triangle (λ(x::real^2). abs (det (matrix ?T)))
    using measure-differentiable-image-eq-affine[OF lmeasurable-S assms(2)] by
auto
  ultimately show ?thesis by auto
qed

```

lemma *triangle-affine-preserves-interior*:
assumes $A = \text{triangle-affine } a \ b \ c$ **and** $L = \text{triangle-linear } a \ b \ c$
assumes $\neg \text{collinear } \{a, b, c\}$
shows $A \text{ ' } (\text{interior } S) = \text{interior } (A \text{ ' } S)$
proof –
let $?trans = \lambda x::\text{real}^2. x + a$
have *linear* L **by** (*simp add: assms(2) triangle-linear-def*)
moreover **have** *surj* L
using *triangle-linear-inj[of a b c] linear-injective-imp-surjective[of L] assms calculation*
by *blast*
ultimately **have** $L: \text{interior}(L \text{ ' } S) = L \text{ ' } (\text{interior } S)$
using *interior-surjective-linear-image* **by** *blast*
moreover **have** $\text{interior} (?trans \text{ ' } S) = ?trans \text{ ' } (\text{interior } S)$
using *interior-translation*
by (*metis (no-types, lifting) add commute image-cong*)
moreover **have** $A = ?trans \circ L$ **using** *assms triangle-affine-def triangle-linear-def*
by *fastforce*
ultimately **show** *?thesis*
by (*smt (verit, del-insts) add commute image-comp image-cong interior-translation*)
qed

lemma *triangle-affine-preserves-affine-hull*:
assumes $A = \text{triangle-affine } a \ b \ c$
assumes $\neg \text{collinear } \{a, b, c\}$
shows $A \text{ ' } (\text{affine hull } S) = \text{affine hull } (A \text{ ' } S)$
proof –
let $?L = \text{triangle-linear } a \ b \ c$
have *linear* $?L$ **by** (*simp add: triangle-linear-def*)
then **have** $?L \text{ ' } (\text{affine hull } S) = \text{affine hull } (?L \text{ ' } S)$
by (*simp add: affine-hull-linear-image linear-linear*)
then **show** *?thesis*
unfolding *assms(1) triangle-affine-def*
by (*metis affine-hull-translation image-image triangle-linear-def*)
qed

lemma *triangle-measure-convex-hull-measure-path-inside-same*:
assumes *p-triangle*: $p = \text{make-triangle } a \ b \ c$
assumes *elem-triangle*: $\text{elem-triangle } a \ b \ c$
shows $\text{measure lebesgue } (\text{convex hull } \{a, b, c\}) = \text{measure lebesgue } (\text{path-inside } p)$
(is $\text{measure lebesgue } ?S = \text{measure lebesgue } ?I$ **)**
proof –
have *bounded* $?S$ **by** (*simp add: finite-imp-bounded-convex-hull*)
then **have** $\text{measure lebesgue } (\text{frontier } ?S) = \text{measure lebesgue } ?S - \text{measure lebesgue } (\text{interior } ?S)$
using *measure-frontier[of ?S]* **by** *auto*
then **have** $\dots = 0$

by (*metis convex-convex-hull negligible-convex-frontier negligible-imp-measure0*)
moreover have $?I = \text{interior } ?S$
using *assms triangle-is-convex*
by (*metis (no-types, lifting) make-triangle-def convex-polygon-inside-is-convex-hull-interior empty-set insert-absorb2 insert-commute list.simps(15) elem-triangle-def triangle-is-polygon*)
ultimately show $?thesis$ **by auto**
qed

lemma *on-triangle-path-image-cases*:

assumes $p = \text{make-triangle } a \ b \ c$
assumes $d \in \text{path-image } p$
shows $d \in \text{path-image } (\text{linepath } a \ b) \vee d \in \text{path-image } (\text{linepath } b \ c) \vee d \in \text{path-image } (\text{linepath } c \ a)$
using *assms unfolding make-triangle-def*
by (*metis make-polygonal-path.simps(3) make-polygonal-path.simps(4) not-in-path-image-join*)

lemma *on-triangle-frontier-cases*:

fixes $a \ b \ c :: \text{real}^2$
assumes $\neg \text{collinear } \{a, b, c\}$
assumes $d \in \text{frontier } (\text{convex hull } \{a, b, c\})$
shows $d \in \text{path-image } (\text{linepath } a \ b) \vee d \in \text{path-image } (\text{linepath } b \ c) \vee d \in \text{path-image } (\text{linepath } c \ a)$

proof –

let $?p = \text{make-triangle } a \ b \ c$
have *polygon* $?p$ **by** (*simp add: assms(1) triangle-is-polygon*)
then have *path-image* $?p = \text{frontier } (\text{convex hull } \{a, b, c\})$
unfolding *make-triangle-def*
by (*smt (verit, ccfv-threshold) assms(1) convex-polygon-frontier-is-path-image2 convex-polygon-is-convex-hull empty-set insert-absorb2 insert-commute list.simps(15) make-triangle-def polygon-convex-iff sup-commute triangle-is-convex*)
thus $?thesis$ **using** *on-triangle-path-image-cases assms(2)* **by blast**
qed

lemma *triangle-path-image-subset-convex*:

assumes $p = \text{make-triangle } a \ b \ c$
shows $\text{path-image } p \subseteq \text{convex hull } \{a, b, c\}$
using *polygon-path-image-subset-convex polygon-at-least-3-vertices make-triangle-def*
by (*metis (no-types, lifting) assms empty-set insert-absorb2 insert-commute insert-iff length-pos-if-in-set list.simps(15)*)

lemma *triangle-convex-hull*:

assumes $p = \text{make-triangle } a \ b \ c$ **and** $\neg \text{collinear } \{a, b, c\}$
shows $\text{convex hull } \{a, b, c\} = (\text{path-image } p) \cup (\text{path-inside } p)$
using *triangle-is-convex[OF assms(1) assms(2)]*
by (*smt (z3) Un-commute assms(1) assms(2) closure-Un-frontier convex-closure convex-polygon-is-convex-hull insert-absorb2 insert-commute inside-outside-def inside-outside-polygon list.set(1) list.set(2) make-triangle-def triangle-is-polygon*)

```

end
theory Unit-Geometry
imports
  HOL-Analysis.Polytope
  Polygon-Jordan-Curve
  Triangle-Lemmas

```

```
begin
```

21 Measure Setup

```
lemma finite-convex-is-measurable:
```

```
  fixes  $p :: (\text{real}^2)$  set
```

```
  assumes  $p = \text{convex hull } l$  and finite  $l$ 
```

```
  shows  $p \in \text{sets lebesgue}$ 
```

```
proof -
```

```
  have polytope  $p$ 
```

```
    unfolding polytope-def using assms by force
```

```
  hence compact  $p$  using polytope-imp-compact by auto
```

```
  thus ?thesis using lmeasurable-compact by blast
```

```
qed
```

```
lemma unit-square-lebesgue: unit-square  $\in$  sets lebesgue
```

```
  using finite-convex-is-measurable by auto
```

```
lemma unit-triangle-lebesgue: unit-triangle  $\in$  sets lebesgue
```

```
  using finite-convex-is-measurable by auto
```

```
lemma unit-triangle-lmeasurable: unit-triangle  $\in$  lmeasurable
```

```
  by (simp add: bounded-convex-hull bounded-set-imp-lmeasurable unit-triangle-lebesgue)
```

22 Unit Triangle

```
lemma unit-triangle-vts-not-collinear:
```

```
   $\neg$  collinear  $\{(\text{vector } [0, 0])::\text{real}^2, \text{vector } [1, 0], \text{vector } [0, 1]\}$ 
```

```
  (is  $\neg$  collinear  $\{?a, ?b, ?c\}$ )
```

```
proof(rule ccontr)
```

```
  assume  $\neg \neg$  collinear  $\{?a, ?b, ?c\}$ 
```

```
  then have collinear  $\{?a, ?b, ?c\}$  by auto
```

```
  then obtain  $u :: \text{real}^2$  where  $u: u \neq 0 \wedge$ 
```

```
     $(\forall x \in \{?a, ?b, ?c\}. \forall y \in \{?a, ?b, ?c\}. \exists c. x - y = c *_R u)$ 
```

```
  by (meson collinear)
```

```
  then obtain  $c1\ c2$  where  $c1: ?b - ?a = c1 *_R u$  and  $c2: ?c - ?a = c2 *_R u$ 
```

```
by blast
```

```
  then have  $c1 *_R u = ?b$ 
```

```
  by (metis (no-types, opaque-lifting) diff-zero scaleR-eq-0-iff vector-2(1) vector-2(2) vector-minus-component vector-scaleR-component zero-neq-one)
```

```
  moreover have  $c2 *_R u = ?c$  using  $c1\ c2$  calculation by force
```

ultimately have $u\$1 = 0 \wedge u\$2 = 0$
by (*metis scaleR-eq-0-iff vector-2(1) vector-2(2) vector-scaleR-component zero-neq-one*)
then have $u = 0$
by (*metis (mono-tags, opaque-lifting) exhaust-2 vec-eq-iff zero-index*)
moreover have $u \neq 0$ **using** u **by** *auto*
ultimately show *False* **by** *auto*
qed

lemma *unit-triangle-convex*:

assumes $p = (\text{make-polygonal-path } [\text{vector } [0, 0], \text{vector } [1, 0], \text{vector } [0, 1], \text{vector } [0, 0]])$
(is $p = \text{make-polygonal-path } [?O, ?e1, ?e2, ?O]$ **)**
shows *convex (path-inside p)*
proof –
have $\neg \text{collinear } \{?O, ?e1, ?e2\}$ **by** (*simp add: unit-triangle-pts-not-collinear*)
thus *?thesis* **using** *triangle-is-convex make-triangle-def assms* **by** *force*
qed

lemma *unit-triangle-char*:

shows $\text{unit-triangle} = \{x. 0 \leq x \$ 1 \wedge 0 \leq x \$ 2 \wedge x \$ 1 + x \$ 2 \leq 1\}$
(is $\text{unit-triangle} = ?S$ **)**

proof –

have $\text{unit-triangle} \subseteq ?S$
proof(*rule subsetI*)
fix x **assume** $x \in \text{unit-triangle}$
then obtain $a b c$ **where**
 $x = a *_R (\text{vector } [0, 0]) + b *_R (\text{vector } [1, 0]) + c *_R (\text{vector } [0, 1])$
 $\wedge a \geq 0 \wedge b \geq 0 \wedge c \geq 0 \wedge a + b + c = 1$
using *convex-hull-3* **by** *blast*
thus $x \in \{x. 0 \leq x \$ 1 \wedge 0 \leq x \$ 2 \wedge x \$ 1 + x \$ 2 \leq 1\}$ **by** *simp*
qed
moreover have $?S \subseteq \text{unit-triangle}$
proof(*rule subsetI*)
fix x **assume** $x \in ?S$
then obtain $b c$ **where** $bc: x\$1 = b \wedge x\$2 = c \wedge 0 \leq b \wedge 0 \leq c \wedge b + c \leq 1$ **by** *blast*
moreover then obtain a **where** $a \geq 0 \wedge a + b + c = 1$ **using** *that[of 1 - b - c]* **by** *argo*
moreover have $a *_R ((\text{vector } [0, 0])::(\text{real}^2)) = \text{vector } [0, 0]$ **by** (*simp add: vec-scaleR-2*)
moreover have $x = (a *_R \text{vector } [0, 0]) + (b *_R \text{vector } [1, 0]) + (c *_R \text{vector } [0, 1])$
using *segment-horizontal bc* **by** *fastforce*
ultimately show $x \in \text{unit-triangle}$ **using** *convex-hull-3* **by** *blast*
qed
ultimately show *?thesis* **by** *blast*
qed

lemma *unit-triangle-interior-char*:

shows $interior\ unit\ triangle = \{x. 0 < x \$ 1 \wedge 0 < x \$ 2 \wedge x \$ 1 + x \$ 2 < 1\}$
 (is $interior\ unit\ triangle = ?S$)
proof –
 have $interior\ unit\ triangle \subseteq ?S$
proof(rule subsetI)
 fix x **assume** $x \in interior\ unit\ triangle$
moreover have $DIM(real^2) = 2$ **by** simp
ultimately obtain $a\ b\ c$ **where**
 $x = a *_R (vector\ [0,\ 0]) + b *_R (vector\ [1,\ 0]) + c *_R (vector\ [0,\ 1])$
 $\wedge a > 0 \wedge b > 0 \wedge c > 0 \wedge a + b + c = 1$
using $interior\ convex\ hull\ 3\ minimal$ [of $(vector\ [0,\ 0])::(real^2)$ $(vector\ [1,\ 0])::(real^2)$ $(vector\ [0,\ 1])::(real^2)$]
using $unit\ triangle\ vts\ not\ collinear$
by auto
thus $x \in \{x. 0 < x \$ 1 \wedge 0 < x \$ 2 \wedge x \$ 1 + x \$ 2 < 1\}$ **by** simp
qed
moreover have $?S \subseteq interior\ unit\ triangle$
proof(rule subsetI)
 fix x **assume** $x \in ?S$
then obtain $b\ c$ **where** $bc: x\$1 = b \wedge x\$2 = c \wedge 0 < b \wedge 0 < c \wedge b + c < 1$ **by** blast
moreover then obtain a **where** $a > 0 \wedge a + b + c = 1$ **using** that[of $1 - b - c$] **by** argo
moreover have $a *_R ((vector\ [0,\ 0])::(real^2)) = vector\ [0,\ 0]$ **by** (simp add: vec-scaleR-2)
moreover have $x = (a *_R vector\ [0,\ 0]) + (b *_R vector\ [1,\ 0]) + (c *_R vector\ [0,\ 1])$
using $segment\ horizontal\ bc$ **by** fastforce
moreover have $DIM(real^2) = 2$ **by** simp
ultimately show $x \in interior\ unit\ triangle$
using $interior\ convex\ hull\ 3\ minimal$ [of $(vector\ [0,\ 0])::(real^2)$ $(vector\ [1,\ 0])::(real^2)$ $(vector\ [0,\ 1])::(real^2)$]
using $unit\ triangle\ vts\ not\ collinear$
by fast
qed
ultimately show $?thesis$ **by** blast
qed

lemma $unit\ triangle\ is\ elementary: elem\ triangle\ (vector\ [0,\ 0])\ (vector\ [1,\ 0])\ (vector\ [0,\ 1])$
 (is $elem\ triangle\ ?a\ ?b\ ?c$)
proof –
let $?UT = unit\ triangle$
have $\neg collinear\ \{?a,\ ?b,\ ?c\}$ **using** $unit\ triangle\ vts\ not\ collinear$ **by** auto
moreover have $integral\ vec\ ?a \wedge integral\ vec\ ?b \wedge integral\ vec\ ?c$
by (simp add: integral-vec-def is-int-def)
moreover have $\{x \in ?UT. integral\ vec\ x\} = \{?a,\ ?b,\ ?c\}$ (is $?UT\ integral = ?abc$)

```

proof –
  have ?UT-integral  $\supseteq$  ?abc using calculation(2) hull-subset by fastforce
  moreover have ?UT-integral  $\subseteq$  ?abc
  proof –
    have  $\bigwedge x. x \in \text{unit-triangle} \implies \text{integral-vec } x \implies x \neq \text{vector } [0, 0] \implies x \neq$ 
     $\text{vector } [1, 0] \implies x \neq \text{vector } [0, 1] \implies \text{False}$ 
    proof –
      fix x
      assume *:  $x \in \text{unit-triangle}$ 
        integral-vec x
         $x \neq \text{vector } [0, 0]$ 
         $x \neq \text{vector } [1, 0]$ 
         $x \neq \text{vector } [0, 1]$ 
      then have x-inset:  $x \in \{x. 0 \leq x \$ 1 \wedge 0 \leq x \$ 2 \wedge x \$ 1 + x \$ 2 \leq 1\}$ 
        using unit-triangle-char by auto
      have  $x \$ 1 = 1 \implies x \$ 2 \neq 0$ 
        using *
        by (smt (verit, del-insts) exhaust-2 vec-eq-iff vector-2(1) vector-2(2))
      then have  $x \$ 1 = 1 \implies x \$ 1 + x \$ 2 > 1 \vee x \$ 2 < 0$ 
        using *(2) unfolding integral-vec-def is-int-def
        by linarith
      then have x1-not-1:  $x \$ 1 = 1 \implies \text{False}$ 
        using x-inset by simp
      have  $x \$ 1 = 0 \implies x \$ 2 \neq 0 \wedge x \$ 2 \neq 1$ 
        using *
        by (smt (verit, del-insts) exhaust-2 vec-eq-iff vector-2(1) vector-2(2))
      then have  $x \$ 1 = 0 \implies x \$ 1 + x \$ 2 > 1 \vee x \$ 1 + x \$ 2 < 0$ 
        using *(2) unfolding integral-vec-def is-int-def
        by auto
      then have x1-not-0:  $x \$ 1 = 0 \implies \text{False}$ 
        using x-inset by simp
      have x1-not-lt0:  $x \$ 1 < 0 \implies \text{False}$ 
        using x-inset by auto
      have x1-not-gt1:  $x \$ 1 > 1 \implies \text{False}$ 
        using x-inset by auto
      then show False using x1-not-0 x1-not-1 x1-not-lt0 x1-not-gt1
        using *(2) unfolding integral-vec-def is-int-def
        by force
    qed
  then have  $\exists x \in ?UT\text{-integral}. x \notin ?abc \wedge \text{integral-vec } x \implies \text{False}$ 
    by blast
  then show ?thesis by blast
qed
ultimately show ?thesis by blast
qed
ultimately show ?thesis unfolding elem-triangle-def by auto
qed

```

lemma *unit-triangles-same-area*:

measure lebesgue unit-triangle' = measure lebesgue unit-triangle
proof –
let ?a = (vector [1, 1])::real^2
let ?b = (vector [0, 1])::real^2
let ?c = (vector [1, 0])::real^2
let ?A = triangle-affine ?a ?b ?c
let ?L = triangle-linear ?a ?b ?c
have collinear-second-component: $\bigwedge c::\text{real}^2. \text{collinear } \{?a, ?b, c\} \implies c \ \$ \ 2 =$
1
proof –
fix p
assume collinear {?a, ?b, p}
then obtain u **where** u-prop: $\forall x \in \{\text{vector } [1, 1], \text{vector } [0, 1], p\}.$
 $\forall y \in \{\text{vector } [1, 1], \text{vector } [0, 1], p\}. \exists c. x - y = c *_R u$
unfolding collinear-def **by** auto
then have c-ab: $\exists c. ?a - ?b = c *_R u$
by blast
then have u-2: $u \ \$ \ 2 = 0$
using vector-2
by (metis cancel-comm-monoid-add-class.diff-cancel diff-zero scaleR-eq-0-iff
vector-minus-component vector-scaleR-component zero-neq-one)
have u-1: $u \ \$ \ 1 \neq 0$
using c-ab vector-2
by (smt (z3) scaleR-right-diff-distrib vector-minus-component vector-scaleR-component)
then have $(\exists c. ?a - p = c *_R u) \wedge (\exists c. ?b - p = c *_R u)$
using u-prop **by** blast
then show $p \ \$ \ 2 = 1$
using u-1 u-2
by (metis eq-iff-diff-eq-0 scaleR-zero-right vector-2(2) vector-minus-component
vector-scaleR-component)
qed
have unit-triangle' = convex hull {?a, ?b, ?c} **by** (simp add: insert-commute)
then have ?A ' unit-triangle = unit-triangle' **using** triangle-affine-img[of ?a ?b
?c] **by** argo
moreover have abs (det (matrix ?L)) = 1
proof –
have matrix ?L = transpose (vector [?b - ?a, ?c - ?a])
unfolding triangle-linear-def
by (simp add: triangle-mat-def)
also have det ... = det (vector [?b - ?a, ?c - ?a]) **using** det-transpose **by**
blast
also have ... = $(?b - ?a) \ \$ \ 1 * (?c - ?a) \ \$ \ 2 - (?c - ?a) \ \$ \ 1 * (?b - ?a) \ \$ \ 2$
using det-2 **by** (metis mult.commute vector-2(1) vector-2(2))
finally show ?thesis **by** simp
qed
moreover have $\neg \text{collinear } \{?a, ?b, ?c\}$ **using** collinear-second-component vec-
tor-2 **by** force
ultimately have *measure lebesgue unit-triangle' = integral unit-triangle* ($\lambda(x::\text{real}^2).$
1)

```

    using triangle-measure-integral-of-det[of ?a ?b ?c]
    by (smt (verit, ccfv-SIG) Henstock-Kurzweil-Integration.integral-cong insert-commute)
    also have ... = measure lebesgue unit-triangle
    by (simp add: lmeasure-integral unit-triangle-lmeasurable)
    finally show ?thesis .
qed

```

23 Unit Square

lemma *convex-hull-4*:

$\text{convex hull } \{a,b,c,d\} = \{ u *_R a + v *_R b + w *_R c + t *_R d \mid u v w t. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge 0 \leq t \wedge u + v + w + t = 1 \}$

proof –

have *fin*: *finite* $\{a,b,c,d\}$ *finite* $\{b,c,d\}$ *finite* $\{c,d\}$ *finite* $\{d\}$

by *auto*

have *: $\bigwedge x y z w :: \text{real}. x + y + z + w = 1 \longleftrightarrow x = 1 - y - z - w$

by (*auto simp: field-simps*)

show *?thesis*

unfolding *convex-hull-finite*[*OF fin(1)*]

unfolding *convex-hull-finite-step*[*OF fin(2)*]

unfolding *convex-hull-finite-step*[*OF fin(3)*]

unfolding *convex-hull-finite-step*[*OF fin(4)*]

unfolding *

apply *auto*

apply (*smt (verit, ccfv-threshold) add.commute diff-add-cancel diff-diff-eq*)

subgoal for *v w t*

apply (*rule exI* [**where** $x=1 - v - w - t$], *simp*)

apply (*rule exI* [**where** $x=v$], *simp*)

apply (*rule exI* [**where** $x=w$], *simp*)

apply (*rule exI* [**where** $x=\lambda x. t$], *simp*)

done

done

qed

lemma *unit-square-characterization-helper*:

fixes *a b* :: *real*

assumes $0 \leq a \wedge a \leq 1 \wedge 0 \leq b \wedge b \leq 1$ **and**

$a \leq b$

obtains *u v w t* **where**

$\text{vector } [a, b] = u *_R ((\text{vector } [0, 0])::\text{real}^2)$

$+ v *_R (\text{vector } [0, 1])$

$+ w *_R (\text{vector } [1, 1])$

$+ t *_R (\text{vector } [1, 0])$

$\wedge 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge 0 \leq t \wedge u + v + w + t = 1$

proof–

let *?a* = $(\text{vector } [0, 0])::(\text{real}^2)$

let *?b* = $(\text{vector } [0, 1])::(\text{real}^2)$

let *?c* = $(\text{vector } [1, 1])::(\text{real}^2)$

let *?d* = $(\text{vector } [1, 0])::(\text{real}^2)$

let $?w = a$
let $?v = b - a$
let $?u = (1 - ?w - ?v)::\text{real}$
let $?t = 0::\text{real}$
let $?T = \{u *_R ?a + v *_R ?b + w *_R ?c + t *_R ?d \mid u \ v \ w \ t. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge 0 \leq t \wedge u + v + w + t = 1\}$
have $?u *_R ?a = 0$
by (*smt (verit, del-insts) exhaust-2 scaleR-zero-right vec-eq-iff vector-2(1) vector-2(2) zero-index*)
moreover have $?w *_R ?c = \text{vector } [a, a]$
proof-
have $(?w *_R ?c)\$1 = a$ **by** *simp*
moreover have $(?w *_R ?c)\$2 = a$ **by** *simp*
ultimately show *?thesis* **by** (*smt (verit) vec-eq-iff exhaust-2 vector-2(1) vector-2(2)*)
qed
moreover have $?v *_R ?b = \text{vector } [0, b - a]$
proof-
have $(?v *_R ?b)\$1 = 0$ **by** *fastforce*
moreover have $(?v *_R ?b)\$2 = b - a$ **by** *simp*
ultimately show *?thesis* **by** (*smt (verit) vec-eq-iff exhaust-2 vector-2(1) vector-2(2)*)
qed
ultimately have $?u *_R ?a + ?v *_R ?b + ?w *_R ?c + ?t *_R ?d = \text{vector } [0, b - a] + \text{vector } [a, a]$
by *fastforce*
also have $\dots = \text{vector } [a, b]$
by (*smt (verit, del-insts) diff-add-cancel exhaust-2 vec-eq-iff vector-2(1) vector-2(2) vector-add-component*)
finally have $\text{vector } [a, b] = ?u *_R ?a + ?v *_R ?b + ?w *_R ?c + ?t *_R ?d$ **by** *presburger*
moreover have $0 \leq ?u \wedge ?u \leq 1 \wedge 0 \leq ?v \wedge ?v \leq 1$ **using** *assms* **by** *simp*
moreover have $0 \leq ?w \wedge ?w \leq 1 \wedge 0 \leq ?t \wedge ?t \leq 1$ **using** *assms* **by** *simp*
moreover have $?u + ?v + ?w + ?t = 1$ **by** *argo*
ultimately show *?thesis* **using** *that[of ?u ?v ?w ?t]* **by** *blast*
qed

lemma *unit-square-characterization:*

unit-square = $\{x. 0 \leq x\$1 \wedge x\$1 \leq 1 \wedge 0 \leq x\$2 \wedge x\$2 \leq 1\}$ (**is** *unit-square* = *?S*)

proof-

let $?a = (\text{vector } [0, 0])::(\text{real}^2)$
let $?b = (\text{vector } [0, 1])::(\text{real}^2)$
let $?c = (\text{vector } [1, 1])::(\text{real}^2)$
let $?d = (\text{vector } [1, 0])::(\text{real}^2)$
let $?T = \{u *_R ?a + v *_R ?b + w *_R ?c + t *_R ?d \mid u \ v \ w \ t. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge 0 \leq t \wedge u + v + w + t = 1\}$
have *unit-square* = $?T$ **using** *convex-hull-4* **by** *blast*
moreover have $?T \subseteq ?S$

```

proof(rule subsetI)
  fix x
  assume x ∈ ?T
  then obtain u v w t where x = u *R ?a + v *R ?b + w *R ?c + t *R ?d and
    0 ≤ u and 0 ≤ v and 0 ≤ w and 0 ≤ t and u + v + w + t = 1 by auto
  moreover from this have
    x$1 = u * 0 + v * 0 + w * 1 + t * 1 ∧ x$2 = u * 0 + v * 1 + w * 1 +
t * 0 by simp
  ultimately have 0 ≤ x$1 ∧ x$1 ≤ 1 ∧ 0 ≤ x$2 ∧ x$2 ≤ 1 by linarith
  thus x ∈ ?S by blast
qed
moreover have ?S ⊆ ?T
proof(rule subsetI)
  fix x :: real^2
  assume *: x ∈ ?S
  { assume x$1 < x$2
    then have x$1 ≤ x$2 by fastforce
    then obtain u v w t where vector [x$1, x$2] = u *R ?a + v *R ?b + w *R
?c + t *R ?d ∧ 0 ≤ u ∧ 0 ≤ v ∧ 0 ≤ w ∧ 0 ≤ t ∧ u + v + w + t = 1
    using * unit-square-characterization-helper[of x$1 x$2] by blast
    moreover have x = vector [x$1, x$2]
    by (smt (verit, ccfv-threshold) exhaust-2 vec-eq-iff vector-2(1) vector-2(2))
    ultimately have x ∈ ?T by force
  } moreover
  { assume x$1 ≥ x$2
    then obtain u v w t where **: vector [x$2, x$1] = u *R ?a + v *R ?b +
w *R ?c + t *R ?d ∧ 0 ≤ u ∧ 0 ≤ v ∧ 0 ≤ w ∧ 0 ≤ t ∧ u + v + w + t = 1
    using * unit-square-characterization-helper[of x$2 x$1] by blast
    have x1: x$1 = v + w using **
    by (smt (verit, ccfv-threshold) mult-cancel-left1 real-scaleR-def scaleR-zero-right
vector-2(2) vector-add-component vector-scaleR-component)
    have x2: x$2 = w + t using **
    by (smt (verit) mult-cancel-left1 real-scaleR-def scaleR-zero-right vector-2(1)
vector-add-component vector-scaleR-component)
    have (u *R ?a + t *R ?b + w *R ?c + v *R ?d)$1 = w + v by auto
    moreover have (u *R ?a + t *R ?b + w *R ?c + v *R ?d)$2 = t + w by
fastforce
    ultimately have u *R ?a + t *R ?b + w *R ?c + v *R ?d = vector [w +
v, t + w]
    by (smt (verit) vec-eq-iff exhaust-2 vector-2(1) vector-2(2))
    also have ... = x using x1 x2
    by (smt (verit, del-insts) add commute exhaust-2 vec-eq-iff vector-2(1)
vector-2(2))
    ultimately have x ∈ ?T
    by (smt (verit, ccfv-SIG) ** mem-Collect-eq)
  }
ultimately show x ∈ ?T by argo
qed
ultimately show ?thesis by auto

```

qed

lemma *e1e2-basis*:

defines $e1 \equiv (\text{vector } [1, 0])::(\text{real}^2)$ and

$e2 \equiv (\text{vector } [0, 1])::(\text{real}^2)$

shows $e1 = \text{axis } 1 (1::\text{real})$ and $e1 \in (\text{Basis}::((\text{real}^2) \text{ set}))$ and

$e2 = \text{axis } 2 (1::\text{real})$ and $e2 \in (\text{Basis}::((\text{real}^2) \text{ set}))$

proof –

have $(1::\text{real}) \in \text{Basis}$ by *simp*

then have $\text{axis } 1 (1::\text{real}) \in (\bigcup i. \bigcup u \in (\text{Basis}::(\text{real set})). \{\text{axis } i u\})$ by *blast*

moreover show $e1\text{-axis}: e1 = \text{axis } 1 (1::\text{real})$

unfolding *axis-def vector-def e1-def* by *auto*

ultimately show $e1\text{-basis}: e1 \in (\text{Basis}::((\text{real}^2) \text{ set}))$ by *simp*

have $(1::\text{real}) \in \text{Basis}$ by *simp*

then have $\text{axis } 2 (1::\text{real}) \in (\bigcup i. \bigcup u \in (\text{Basis}::(\text{real set})). \{\text{axis } i u\})$ by *blast*

moreover show $e2\text{-axis}: e2 = \text{axis } 2 (1::\text{real})$

unfolding *axis-def vector-def e2-def* by *auto*

ultimately show $e2\text{-basis}: e2 \in (\text{Basis}::((\text{real}^2) \text{ set}))$ by *simp*

qed

lemma *unit-square-cbox*: $\text{unit-square} = \text{cbox } (\text{vector } [0, 0]) (\text{vector } [1, 1])$

proof –

let $?O = (\text{vector } [0, 0])::(\text{real}^2)$

let $?e1 = (\text{vector } [1, 0])::(\text{real}^2)$

let $?e2 = (\text{vector } [0, 1])::(\text{real}^2)$

let $?I = (\text{vector } [1, 1])::(\text{real}^2)$

let $?cbox = \{x. \forall i \in \text{Basis}. ?O \cdot i \leq x \cdot i \wedge x \cdot i \leq ?I \cdot i\}$

have $\text{unit-square} = \{x. 0 \leq x\$1 \wedge x\$1 \leq 1 \wedge 0 \leq x\$2 \wedge x\$2 \leq 1\}$ (is $\text{unit-square} = ?S$)

using *unit-square-characterization* by *auto*

moreover have $?S \subseteq ?cbox$

proof (rule *subsetI*)

fix x

assume *: $x \in ?S$

have $?O \cdot ?e1 \leq x \cdot ?e1 \wedge x \cdot ?e1 \leq ?I \cdot ?e1$

using *e1e2-basis*

by (smt (verit, del-Insts) * *cart-eq-inner-axis mem-Collect-eq vector-2(1)*)

moreover have $?O \cdot ?e2 \leq x \cdot ?e2 \wedge x \cdot ?e2 \leq ?I \cdot ?e2$

using *e1e2-basis*

by (smt (verit, del-Insts) * *cart-eq-inner-axis mem-Collect-eq vector-2(2)*)

ultimately show $x \in ?cbox$

by (smt (verit, best) * *axis-index cart-eq-inner-axis exhaust-2 mem-Collect-eq vector-2(1) vector-2(2)*)

qed

moreover have $?cbox \subseteq ?S$

proof (rule *subsetI*)

fix $x :: \text{real}^2$

```

assume *:  $x \in ?cbox$ 
then have  $0 \leq ?e1 \cdot x$  using  $e1e2$ -basis
  by (metis (no-types, lifting) cart-eq-inner-axis inner-commute mem-Collect-eq
vector-2(1))
  moreover have  $?e1 \cdot x \leq 1$  using  $e1e2$ -basis
  by (smt (verit, ccfv-SIG) * inner-axis inner-commute mem-Collect-eq real-inner-1-right
vector-2(1))
  moreover have  $0 \leq ?e2 \cdot x$ 
  by (metis (no-types, lifting) * cart-eq-inner-axis e1e2-basis(3) e1e2-basis(4)
inner-commute mem-Collect-eq vector-2(2))
  moreover have  $?e2 \cdot x \leq 1$ 
  by (metis (no-types, lifting) * cart-eq-inner-axis e1e2-basis(3) e1e2-basis(4)
inner-commute mem-Collect-eq vector-2(2))
  moreover have  $?e1 \cdot x = x\$1$ 
  by (simp add: cart-eq-inner-axis e1e2-basis inner-commute)
  moreover have  $?e2 \cdot x = x\$2$ 
  by (simp add: cart-eq-inner-axis e1e2-basis inner-commute)
  ultimately show  $x \in ?S$  by force
qed
ultimately show ?thesis unfolding cbox-def by order
qed

```

lemma *unit-square-area: measure lebesgue unit-square = 1*

proof–

```

let  $?e1 = (\text{vector } [1, 0]) :: (\text{real}^2)$ 
let  $?e2 = (\text{vector } [0, 1]) :: (\text{real}^2)$ 
have unit-square = cbox (vector [0, 0]) (vector [1, 1]) (is unit-square = cbox
?O ?I)
  using unit-square-cbox by blast
also have emeasure lborel ... = 1 using emeasure-lborel-cbox-eq
proof–
  have  $?I \cdot ?e1 = (1 :: \text{real})$ 
  by (simp add: e1e2-basis(1) inner-axis' inner-commute)
  moreover have  $?I \cdot ?e2 = (1 :: \text{real})$  by (simp add: e1e2-basis(3) inner-axis'
inner-commute)
  ultimately have basis-dot:  $\forall b \in \text{Basis}. ?I \cdot b = 1$ 
  by (metis (full-types) axis-inverse e1e2-basis(1) e1e2-basis(3) exhaust-2)

  have  $?O \cdot ?e1 \leq ?I \cdot ?e1$  by (simp add: e1e2-basis(1) inner-axis)
  moreover have  $?O \cdot ?e2 \leq ?I \cdot ?e2$  by (simp add: e1e2-basis(3) inner-axis)
  ultimately have  $\forall b \in \text{Basis}. ?O \cdot b \leq ?I \cdot b$ 
  by (smt (verit, ccfv-threshold) axis-index cart-eq-inner-axis exhaust-2 insert-iff
vector-2(1) vector-2(2))
  then have emeasure lborel (cbox ?O ?I) =  $(\prod_{b \in \text{Basis}} (?I - ?O) \cdot b)$ 
  using emeasure-lborel-cbox-eq by auto
  also have  $\dots = (\prod_{b \in \text{Basis}} ?I \cdot b)$ 
  by (smt (verit, del-insts) axis-index diff-zero euclidean-all-zero-iff exhaust-2
inner-axis real-inner-1-right vector-2(1) vector-2(2))
  also have  $\dots = (\prod_{b \in \text{Basis}} (1 :: \text{real}))$  using basis-dot by fastforce

```

finally show *?thesis* **by** *simp*
qed
finally have *emeasure lborel unit-square = 1* .
moreover have *emeasure lborel unit-square = measure lebesgue unit-square*
by (*simp add: emeasure-eq-measure2 unit-square-cbox*)
ultimately show *?thesis* **by** *fastforce*
qed

24 Unit Triangle Area is 1/2

lemma *unit-triangle'-char:*

shows *unit-triangle' = {x. x \$ 1 ≤ 1 ∧ x \$ 2 ≤ 1 ∧ x \$ 1 + x \$ 2 ≥ 1}*

proof –

let *?I = (vector [1, 1])::real^2*

let *?e1 = (vector [1, 0])::real^2*

let *?e2 = (vector [0, 1])::real^2*

have *unit-triangle' = {u *_R ?I + v *_R ?e1 + w *_R ?e2 | u v w. 0 ≤ u ∧ 0 ≤ v ∧ 0 ≤ w ∧ u + v + w = 1}*

using *convex-hull-3[of ?I ?e1 ?e2]* **by** *auto*

moreover have $\bigwedge u v w. u *_{\mathbb{R}} ?I + v *_{\mathbb{R}} ?e1 + w *_{\mathbb{R}} ?e2 = ((\text{vector } [u + v, u + w])::\text{real}^2)$

proof –

fix *u v w :: real*

let *?v-e1 = ((vector [v, 0])::real^2)*

let *?w-e2 = ((vector [0, w])::real^2)*

let *?u-I = ((vector [u, u])::real^2)*

have $u *_{\mathbb{R}} ?I = ?u-I$ **using** *vec-scaleR-2* **by** *simp*

moreover have $v *_{\mathbb{R}} ?e1 = ?v-e1$ **using** *vec-scaleR-2* **by** *simp*

moreover have $w *_{\mathbb{R}} ?e2 = ?w-e2$ **using** *vec-scaleR-2* **by** *simp*

ultimately have $1: u *_{\mathbb{R}} ?I + v *_{\mathbb{R}} ?e1 + w *_{\mathbb{R}} ?e2 = ?u-I + ?v-e1 + ?w-e2$

by *argo*

moreover have $(?u-I + ?v-e1 + ?w-e2)$1 = u + v$

using *vector-add-component* **by** *simp*

moreover have $(?u-I + ?v-e1 + ?w-e2)$2 = u + w$

using *vector-add-component* **by** *simp*

ultimately have $?u-I + ?v-e1 + ?w-e2 = ((\text{vector } [u + v, u + w])::\text{real}^2)$

using *vector-2 exhaust-2* **by** (*smt (verit, del-insts) vec-eq-iff*)

thus $u *_{\mathbb{R}} ?I + v *_{\mathbb{R}} ?e1 + w *_{\mathbb{R}} ?e2 = ((\text{vector } [u + v, u + w])::\text{real}^2)$

using *1* **by** *argo*

qed

ultimately have $1: \text{unit-triangle}' = \{(\text{vector } [u + v, u + w])::\text{real}^2 \mid u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1\}$

(**is** *unit-triangle' = ?S*)

by *presburger*

have $\text{unit-triangle}' = \{(\text{vector } [x, y])::\text{real}^2 \mid x y. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1\}$

(**is** *unit-triangle' = ?T*)

proof-
have $\bigwedge x y :: \text{real}. \exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge x = u + v \wedge y = u + w$
 $\implies 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1$ **by force**
moreover have $*$: $\bigwedge x y :: \text{real}. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1$
 $\implies \exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge x = u + v \wedge y = u + w$
proof-
fix $x y :: \text{real}$
let $?u = y + x - 1$
let $?v = 1 - y$
let $?w = 1 - x$
assume $0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge 1 \leq x + y$
then have $0 \leq ?u \wedge 0 \leq ?v \wedge 0 \leq ?w \wedge ?u + ?v + ?w = 1 \wedge x = ?u + ?v \wedge y = ?u + ?w$ **by argo**
thus $\exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge x = u + v \wedge y = u + w$ **by blast**
qed
ultimately have $\forall x y :: \text{real}. ((\exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge x = u + v \wedge y = u + w)$
 $\longleftrightarrow (0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1))$
by metis
then have $\forall z :: \text{real}^2. ((\exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge z\$1 = u + v \wedge z\$2 = u + w)$
 $\longleftrightarrow (0 \leq z\$1 \wedge z\$1 \leq 1 \wedge 0 \leq z\$2 \wedge z\$2 \leq 1 \wedge z\$1 + z\$2 \geq 1))$ **by presburger**
then have $\forall z :: \text{real}^2. ((\exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge z = \text{vector}[u + v, u + w])$
 $\longleftrightarrow (\exists x y. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1 \wedge z = \text{vector}[x, y]))$
by (smt (verit) *)
moreover have $\forall z :: \text{real}^2. z \in ?S \longleftrightarrow (\exists u v w. 0 \leq u \wedge 0 \leq v \wedge 0 \leq w \wedge u + v + w = 1 \wedge z = \text{vector}[u + v, u + w])$
by blast
moreover have $\forall z :: \text{real}^2. z \in ?T \longleftrightarrow (\exists x y. 0 \leq x \wedge x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1 \wedge z = \text{vector}[x, y])$
by blast
ultimately have $?S = ?T$ **by auto**
then show $?thesis$ **using 1 by auto**
qed
moreover have $\{x. 0 \leq x\$1 \wedge x\$1 \leq 1 \wedge 0 \leq x\$2 \wedge x\$2 \leq 1 \wedge x\$1 + x\$2 \geq 1\} \subseteq ?T$
proof(rule subsetI)
fix $z :: \text{real}^2$
assume $*$: $z \in \{x. 0 \leq x\$1 \wedge x\$1 \leq 1 \wedge 0 \leq x\$2 \wedge x\$2 \leq 1 \wedge x\$1 + x\$2 \geq 1\}$
then obtain $x y :: \text{real}$ **where** $z = \text{vector}[x, y] \wedge 0 \leq x$ **using forall-vector-2**
by fastforce
moreover from this have $x \leq 1 \wedge 0 \leq y \wedge y \leq 1 \wedge x + y \geq 1$ **using ***

```

vector-2[of x y] by simp
  ultimately show  $z \in ?T$  by blast
qed
moreover have  $?T \subseteq \{x. 0 \leq x\$1 \wedge x\$1 \leq 1 \wedge 0 \leq x\$2 \wedge x\$2 \leq 1 \wedge x\$1 + x\$2 \geq 1\}$ 
  using vector-2 by force
  ultimately show ?thesis
  by (smt (verit, best) Collect-cong subset-antisym)
qed

lemma unit-square-split-diag:
  shows unit-square = unit-triangle  $\cup$  unit-triangle'
proof -
  let ?S = ({vector [0, 0], vector [0, 1], vector [1, 0]}::(real^2) set)
  let ?S' = ({vector [1, 1], vector [0, 1], vector [1, 0]}::(real^2) set)
  have unit-triangle  $\cup$  unit-triangle'  $\subseteq$  convex hull (?S  $\cup$  ?S') by (simp add: hull-mono)
  moreover have convex hull (?S  $\cup$  ?S')  $\subseteq$  unit-triangle  $\cup$  unit-triangle'
  by (smt (z3) Un-commute Un-left-commute Un-upper1 in-mono insert-is-Un mem-Collect-eq subsetI sup.idem unit-square-characterization unit-triangle-char unit-triangle'-char)
  moreover have unit-square = convex hull (?S  $\cup$  ?S') by (simp add: insert-commute)
  ultimately show ?thesis by blast
qed

lemma unit-triangle-INT-unit-triangle'-measure:
  measure lebesgue (unit-triangle  $\cap$  unit-triangle') = 0
proof -
  let ?e1 = (vector [1, 0])::real^2
  let ?e2 = (vector [0, 1])::real^2
  have unit-triangle  $\cap$  unit-triangle' = { $x::(real^2). 0 \leq x \$ 1 \wedge x \$ 1 \leq 1 \wedge 0 \leq x \$ 2 \wedge x \$ 2 \leq 1 \wedge x \$ 1 + x \$ 2 = 1$ }
  (is unit-triangle  $\cap$  unit-triangle' = ?S)
  using unit-triangle-char unit-triangle'-char
  by auto
  also have ... = path-image (linepath ?e2 ?e1)
  (is ... = ?p)
proof -
  have ?S  $\subseteq$  ?p
  proof (rule subsetI)
    fix  $x :: real^2$ 
    assume  $x \in ?S$ 
    then have *:  $0 \leq 1 - x\$2 \wedge x\$2 = 1 - x\$1 \wedge 0 \leq x\$2 \wedge x\$2 \leq 1$  by simp

    have  $x\$2 *_R ?e2 + x\$1 *_R ?e1 = vector[x\$1, x\$2]$ 
  proof -
    have  $(x\$1 *_R ?e1)\$1 = x\$1$  by simp
    moreover have  $(x\$1 *_R ?e1)\$2 = 0$  by auto
    moreover have  $(x\$2 *_R ?e2)\$1 = 0$  by auto

```

moreover have $(x\$2 *_R ?e2)\$2 = x\$2$ **by** *fastforce*
ultimately have $x\$1 *_R ?e1 = \text{vector } [x\$1, 0] \wedge x\$2 *_R ?e2 = \text{vector } [0, x\$2]$
by *smt (verit, ccfv-SIG) exhaust-2 vec-eq-iff vector-2(1) vector-2(2)*
then have $x\$1 *_R ?e1 + x\$2 *_R ?e2 = \text{vector } [x\$1, 0] + \text{vector } [0, x\$2]$
by *auto*
moreover from this have $(x\$1 *_R ?e1 + x\$2 *_R ?e2)\$1 = x\1 **by** *auto*
moreover from calculation have $(x\$1 *_R ?e1 + x\$2 *_R ?e2)\$2 = x\2
by *auto*
ultimately show *?thesis*
by *(smt (verit) add.commute exhaust-2 vec-eq-iff vector-2(1) vector-2(2))*
qed
also have $\dots = x$
by *(smt (verit, best) exhaust-2 vec-eq-iff vector-2(1) vector-2(2))*
finally have $x\$2 *_R ?e2 + x\$1 *_R ?e1 = x$.
then have $x = (\lambda x. (1 - x) *_R ?e2 + x *_R ?e1) (x\$1) \wedge x\$1 \in \{0..1\}$
using $*$ **by** *auto*
thus $x \in ?p$ **unfolding** *path-image-def linepath-def* **by** *fast*
qed
moreover have $?p \subseteq ?S$
proof(*rule subsetI*)
fix x
assume $*$: $x \in ?p$
then obtain t **where** $*$: $x = (1 - t) *_R ?e2 + t *_R ?e1 \wedge t \in \{0..1\}$
unfolding *path-image-def linepath-def* **by** *blast*
moreover from this have $x\$1 = t$ **by** *simp*
moreover from calculation have $x\$2 = 1 - t$ **by** *simp*
moreover from calculation have $0 \leq t \wedge t \leq 1 \wedge 0 \leq 1 - t \wedge 1 - t \leq 1$
by *simp*
ultimately show $x \in ?S$ **by** *simp*
qed
ultimately show *?thesis* **by** *blast*
qed
also have *measure lebesgue ?p = 0* **using** *linepath-has-measure-0* **by** *blast*
finally show *?thesis* .
qed

lemma *unit-triangle-area: measure lebesgue unit-triangle = 1/2*
proof–
let $?μ = \text{measure lebesgue}$
have $?μ \text{ unit-square} = ?μ \text{ unit-triangle} + ?μ \text{ unit-triangle}'$
using *unit-square-split-diag unit-triangle-INT-unit-triangle'-measure*
by *(simp add: finite-imp-bounded-convex-hull measurable-convex measure-Un3)*
thus *?thesis* **using** *unit-triangles-same-area unit-square-area* **by** *simp*
qed

end
theory *Elementary-Triangle-Area*
imports

begin

25 Area of Elementary Triangle is 1/2

lemma *nonint-in-square-imp-nonint-triangle-imp*:

assumes $A = \text{triangle-affine } a \ b \ c$

assumes $x \in \text{unit-square}$

assumes $\neg \text{integral-vec } x$

assumes $\text{integral-vec } (A \ x)$

assumes $\text{elem-triangle } a \ b \ c$

obtains x' **where** $x' \in \text{unit-triangle} \wedge \neg \text{integral-vec } x' \wedge \text{integral-vec } (A \ x')$

proof –

{ **assume** $x \in \text{unit-triangle}$

then have $?thesis$ **using** *assms that by blast*

} **moreover**

{ **assume** $*: x \notin \text{unit-triangle}$

then have $x \notin \{x. 0 \leq x \ \$ \ 1 \wedge 0 \leq x \ \$ \ 2 \wedge x \ \$ \ 1 + x \ \$ \ 2 \leq 1\}$

using *unit-triangle-char by argo*

then have $x2x1\text{-ge-1}: x \ \$ \ 1 + x \ \$ \ 2 > 1$ **using** *assms(2) unit-square-characterization*

by *force*

let $?x'1 = 1 - x \ \$ \ 1$

let $?x'2 = 1 - x \ \$ \ 2$

let $?x' = \text{vector } [?x'1, ?x'2]$

have $?x'1 + ?x'2 \leq 1$ **using** *x2x1-ge-1 by argo*

then have $?x' \in \text{unit-triangle}$

using *unit-triangle-char assms(2) unit-square-characterization by auto*

moreover have $\neg \text{integral-vec } ?x'$

proof –

have $\neg \text{is-int } (x \ \$ \ 1) \vee \neg \text{is-int } (x \ \$ \ 2)$ **using** *assms(3) unfolding integral-vec-def by blast*

then have $\neg \text{is-int } (?x'1) \vee \neg \text{is-int } (?x'2)$

using *is-int-minus*

by (*metis diff-add-cancel is-int-def minus-diff-eq of-int-1 uminus-add-conv-diff*)

thus $?thesis$ **unfolding** *integral-vec-def by auto*

qed

moreover have $\text{integral-vec } (A \ ?x')$

proof –

let $?L = \text{triangle-linear } a \ b \ c$

have $A\text{-comp}: A = (\lambda x. x + a) \circ ?L$ **by** (*simp add: affine-comp-linear-trans assms(1)*)

then have $Lx\text{-int}: \text{integral-vec } (?L \ x)$

by (*smt (verit, del-insts) assms(4) assms(5) comp-apply diff-add-cancel diff-minus-eq-add integral-vec-minus integral-vec-sum elem-triangle-def*)

have $\text{linear } ?L$ **by** (*simp add: triangle-linear-def*)

moreover have $?L \ ?x' = ?L (\text{vector } [1, 1] - x)$

by (*simp add: mat-vec-mult-2 triangle-linear-def*)

```

ultimately have ?L ?x' = ?L (vector [1, 1]) - ?L x by (simp add: linear-diff)
moreover have integral-vec (?L (vector [1, 1]))
proof-
  have ?L (vector [1, 1]) = vector [(b - a)$1 + (c - a)$1, (b - a)$2 + (c
- a)$2]
  unfolding triangle-linear-def triangle-mat-def transpose-def using mat-vec-mult-2
by simp
  also have ... = (b - a) + (c - a)
  by (smt (verit, del-Insts) exhaust-2 vec-eq-iff vector-2(1) vector-2(2)
vector-add-component)
  finally show ?thesis using assms(5) unfolding elem-triangle-def
  by (metis ab-group-add-class.ab-diff-conv-add-uminus integral-vec-minus
integral-vec-sum)
qed
ultimately have integral-vec (?L ?x')
  using Lx-int integral-vec-sum integral-vec-minus by force
then show ?thesis using A-comp assms(5) integral-vec-sum elem-triangle-def
by auto
qed
ultimately have ?thesis using that by blast
}
ultimately show ?thesis by blast
qed

```

lemma *elem-triangle-integral-mat-bij*:

```

fixes a b c :: real^2
assumes elem-triangle a b c
defines L ≡ triangle-mat a b c
shows integral-mat-bij L

```

proof–

```

let ?A = triangle-affine a b c

```

```

have L: L = transpose (vector [b - a, c - a]) (is L = transpose (vector [?w1,
?w2]))

```

```

  unfolding triangle-mat-def L-def by auto

```

```

have integral-vec ?w1 ∧ integral-vec ?w2

```

```

  by (metis ab-group-add-class.ab-diff-conv-add-uminus assms(1) integral-vec-minus
integral-vec-sum elem-triangle-def)

```

```

then have L-int-entries: ∀ i ∈ {1, 2}. ∀ j ∈ {1, 2}. is-int (L$i$j)

```

```

  by (simp add: L-def triangle-mat-def Finite-Cartesian-Product.transpose-def
integral-vec-def)

```

```

have L-integral: integral-mat L unfolding integral-mat-def

```

```

proof(rule allI)

```

```

  fix v :: real^2

```

```

  show integral-vec v → integral-vec (L * v v)

```

```

proof(rule impI)

```

```

  assume v-int-asm: integral-vec v

```

```

let ?Lv = L * v

have ?Lv$1 = L$1$1 * v$1 + L$1$2 * v$2 by (simp add: mat-vec-mult-2)
then have Lv1-int: is-int (?Lv$1)
  using L-int-entries v-int-assm is-int-sum is-int-mult by (simp add: integral-vec-def)

have ?Lv$2 = L$2$1 * v$1 + L$2$2 * v$2 by (simp add: mat-vec-mult-2)
then have Lv2-int: is-int (?Lv$2)
  using L-int-entries v-int-assm is-int-sum is-int-mult by (simp add: integral-vec-def)

show integral-vec (L * v)
  by (simp add: Lv1-int Lv2-int integral-vec-def)
qed
moreover have integral-mat-surj L
  unfolding integral-mat-surj-def
proof(rule allI)
  fix v :: real^2
  show integral-vec v ⟶ (∃ w. integral-vec w ∧ L * v w = v)
  proof(rule impI)
    assume *: integral-vec v
    obtain w :: real^2 where w: L * v w = v
      using triangle-linear-inj assms(1) full-rank-injective full-rank-surjective
      unfolding elem-triangle-def L-def triangle-linear-def surj-def
      by (smt (verit, best) iso-tuple-UNIV-I)
    moreover have integral-vec w
  proof(rule ccontr)
    assume **: ¬ integral-vec w
    let ?w1 = w$1
    let ?w2 = w$2
    let ?w1' = w$1 - (floor (w$1))
    let ?w2' = w$2 - (floor (w$2))
    let ?w' = (vector [?w1', ?w2']):(real^2)
    have ?w1' ∈ {0..1} ∧ ?w2' ∈ {0..1}
      by (metis add.commute add.right-neutral atLeastAtMost-iff floor-correct
        floor-frac frac-def of-int-0 real-of-int-floor-add-one-ge)
    then have ?w' ∈ unit-square using unit-square-characterization by auto
    moreover have ¬ integral-vec ?w'
      by (metis ** eq-iff-diff-eq-0 floor-frac floor-of-int frac-def integral-vec-def
        is-int-def of-int-0 vector-2(1) vector-2(2))
    moreover have integral-vec (?A ?w')
  proof-
    have ?w' = vector [w$1, w$2] - vector [floor (w$1), floor (w$2)]
      (is ?w' = vector [w$1, w$2] - ?floor-w)
    by (smt (verit, del-Insts) exhaust-2 list.simps(8) list.simps(9) vec-eq-iff
      vector-2(1) vector-2(2) vector-minus-component)
    then have ?w' = w - vector [floor (w$1), floor (w$2)]

```

by (smt (verit, del-insts) exhaust-2 vec-eq-iff vector-2(1) vector-2(2)
 vector-minus-component)
 moreover have ?A ?w' = (L *v ?w') + a **unfolding** triangle-affine-def
 L-def by simp
 ultimately have ?A ?w' = v - (L *v ?floor-w) + a
 by (simp add: matrix-vector-mult-diff-distrib w)
 moreover have integral-vec v \wedge integral-vec a \wedge integral-vec (L *v ?floor-w)
 using * assms(1) L-integral integral-mat-integral-vec integral-vec-2
 unfolding elem-triangle-def
 by blast
 ultimately show ?thesis
 by (metis ab-group-add-class.ab-diff-conv-add-uminus integral-vec-minus
 integral-vec-sum)
 qed
 ultimately obtain w'' where w'': w'' \in unit-triangle \wedge \neg integral-vec w''
 \wedge integral-vec (?A w'')
 using nonint-in-square-img-IMP-nonint-triangle-img[of ?A a b c ?w']
 assms(1) by blast
 moreover have ?A w'' \notin {a, b, c}
 proof-
 have inj ?A using assms(1) elem-triangle-def triangle-affine-inj by auto
 moreover have ?A (vector [0, 0]) = a
 by (metis (no-types, opaque-lifting) add commute add-0 mat-vec-mult-2 ma-
 trix-vector-mult-0-right real-scaleR-def scaleR-zero-right triangle-affine-def zero-index)
 moreover have ?A (vector [1, 0]) = b
 unfolding triangle-affine-def triangle-mat-def transpose-def
 by (metis (no-types) Finite-Cartesian-Product.transpose-def add commute
 column-transpose diff-add-cancel e1e2-basis(1) matrix-vector-mult-basis row-def vec-lambda-eta
 vector-2(1))
 moreover have ?A (vector [0, 1]) = c
 proof-
 have (?A (vector [0, 1]))\$1 = c\$1
 by (metis L-def L add commute column-transpose diff-add-cancel
 e1e2-basis(3) matrix-vector-mult-basis row-def triangle-affine-def vec-lambda-eta vec-
 tor-2(2))
 moreover have (?A (vector [0, 1]))\$2 = c\$2
 by (metis add commute column-transpose diff-add-cancel e1e2-basis(3)
 matrix-vector-mult-basis row-def triangle-affine-def triangle-mat-def vec-lambda-eta
 vector-2(2))
 ultimately show ?thesis by (smt (verit, ccfv-SIG) exhaust-2 vec-eq-iff)
 qed
 moreover have w'' \neq vector [0, 0] \wedge w'' \neq vector [0, 1] \wedge w'' \neq vector
 [1, 0]
 using w'' elem-triangle-def unit-triangle-is-elementary by blast
 ultimately show ?thesis by (metis inj-eq insertE singletonD)
 qed
 moreover have ?A ' unit-triangle = convex hull {a, b, c}
 using triangle-affine-img by blast
 ultimately show False using assms unfolding elem-triangle-def by blast

```

    qed
    ultimately show  $\exists w. \text{integral-vec } w \wedge L * v w = v$  by auto
  qed
  qed
  ultimately show ?thesis unfolding integral-mat-bij-def by auto
  qed

lemma elem-triangle-measure-integral-of-1:
  fixes a b c :: real^2
  defines S  $\equiv$  convex hull {a, b, c}
  assumes elem-triangle a b c
  shows measure lebesgue S = integral unit-triangle ( $\lambda x::\text{real}^2. 1$ )
  proof -
    let ?T = triangle-linear a b c
    have integral-mat-bij (matrix ?T) (is integral-mat-bij ?T-mat)
      by (simp add: assms(2) elem-triangle-integral-mat-bij triangle-linear-def)
    then have abs (det ?T-mat) = 1
      using integral-mat-bij-det-pm1 by fastforce
    thus ?thesis
      using S-def assms(2) triangle-measure-integral-of-det elem-triangle-def by force
  qed

lemma elem-triangle-area-is-half:
  fixes a b c :: real^2
  assumes elem-triangle a b c
  defines S  $\equiv$  convex hull {a, b, c}
  shows measure lebesgue S = 1/2 (is ?S-area = 1/2)
  proof -
    have  $\neg$  collinear {a, b, c} using elem-triangle-def assms(1) by blast
    then have measure lebesgue S = integral unit-triangle ( $\lambda x::\text{real}^2. 1$ )
      using S-def assms(1) elem-triangle-measure-integral-of-1 by blast
    also have ... = measure lebesgue unit-triangle
      using unit-triangle-is-elementary elem-triangle-measure-integral-of-1 unit-triangle-area
      by metis
    finally show ?thesis by (simp add: unit-triangle-area)
  qed

end
theory Pick
imports
  Polygon-Splitting
  Elementary-Triangle-Area
begin

```

26 Setup

26.1 Integral Points Cardinality Properties

lemma bounded-finite:

```

fixes  $A :: (\text{real}^2)$  set
assumes bounded A
shows finite {x::(real^2). integral-vec x ∧ x ∈ A} (is finite ?A-int)
proof –
  obtain  $M$  where  $M: \forall x \in A. \text{norm } x \leq M$  using assms bounded-def by (meson bounded-iff)

  let  $?M\text{-bounded-ints} = \{n. n \in \{-M..M\} \wedge \text{is-int } n\}$ 
  let  $?M\text{-bounded-int-vecs} = \{v::(\text{real}^2). v\$1 \in ?M\text{-bounded-ints} \wedge v\$2 \in ?M\text{-bounded-ints}\}$ 

  have  $\forall x::(\text{real}^2). \text{norm } (x\$1) \leq \text{norm } x \wedge (x\$2) \leq \text{norm } x$ 
  by (smt (verit, ccfv-threshold) Finite-Cartesian-Product.norm-nth-le real-norm-def)
  then have  $\forall x \in ?A\text{-int}. \text{norm } (x\$1) \leq M \wedge \text{norm } (x\$2) \leq M$ 
  using  $M$  dual-order.trans Finite-Cartesian-Product.norm-nth-le by blast
  then have  $\forall x \in ?A\text{-int}. x\$1 \in ?M\text{-bounded-ints} \wedge x\$2 \in ?M\text{-bounded-ints}$ 
  using integral-vec-def intervalE by auto
  then have  $\forall x \in ?A\text{-int}. x \in ?M\text{-bounded-int-vecs}$  by blast
  moreover have finite ?M-bounded-int-vecs
  proof –
    obtain  $S :: \text{int set}$  where  $S: S = \{n. \exists m \in ?M\text{-bounded-ints}. n = m\} \wedge (\forall n \in S. \text{norm } n \leq M)$ 
    by (simp add: abs-le-iff)
    then have finite-S: finite S
    by (metis infinite-int-iff-unbounded le-floor-iff linorder-not-less norm-of-int of-int-abs)

  have finite-M-bounded-ints: finite ?M-bounded-ints
  proof –
    let  $?f = \lambda n::\text{real}. \text{THE } m::\text{int}. n = m$ 
    have  $\forall n \in ?M\text{-bounded-ints}. \exists !m::\text{int}. n = m$  using is-int-def by force
    moreover have inj-on ?f ?M-bounded-ints using inj-on-def is-int-def by
force
    moreover have  $?f \text{ ` } ?M\text{-bounded-ints} \subseteq S$  using calculation S subsetI by
auto
    ultimately show ?thesis using finite-imageD finite-S by (simp add: inj-on-finite)
  qed
  show ?thesis
  proof –
    let  $?f = \lambda x::(\text{real}^2). (\text{THE } m::\text{int}. m = x\$1, \text{THE } n::\text{int}. n = x\$2)$ 
    have inj-on ?f ?M-bounded-int-vecs
    unfolding inj-on-def
    proof clarify
      fix  $x y :: \text{real}^2$ 
      assume  $x1\text{-int}: \text{is-int } (x\$1)$ 
      assume  $x2\text{-int}: \text{is-int } (x\$2)$ 
      assume  $y1\text{-int}: \text{is-int } (y\$1)$ 
      assume  $y2\text{-int}: \text{is-int } (y\$2)$ 
      assume  $x1y1\text{-int-eq}: (\text{THE } m. \text{real-of-int } m = x\$1) = (\text{THE } m. \text{real-of-int}$ 

```

$m = y\$1$
assume $x2y2$ -int-eq: $(THE\ n.\ real-of-int\ n = x\$2) = (THE\ n.\ real-of-int\ n = y\$2)$

have $\exists!m.\ m = x\$1$
by *blast*
moreover have $\exists!n.\ n = y\$1$
by *blast*
moreover have $(THE\ m.\ real-of-int\ m = x\$1) = (THE\ m.\ real-of-int\ m = y\$1)$
using $x1y1$ -int-eq **by** *auto*
ultimately have $x1y1: x\$1 = y\1
using $x1$ -int $y1$ -int *is-int-def* **by** *auto*

have $\exists!m.\ m = x\$2$
by *blast*
moreover have $\exists!n.\ n = y\$2$
by *blast*
moreover have $(THE\ m.\ real-of-int\ m = x\$2) = (THE\ m.\ real-of-int\ m = y\$2)$
using $x2y2$ -int-eq **by** *auto*
ultimately have $x2y2: x\$2 = y\2
using $x2$ -int $y2$ -int *is-int-def* **by** *auto*

show $x = y$ **using** $x1y1\ x2y2$
by $(metis\ (no-types,\ lifting)\ exhaust-2\ vec-eq-iff)$
qed

moreover have $?f\ '?\ M$ -bounded-int-vecs $\subseteq S \times S$
proof(*rule subsetI*)
fix mn
assume $mn \in ?f\ '?\ M$ -bounded-int-vecs
then obtain v **where** v :
 $v \in ?M$ -bounded-int-vecs $\wedge ?f\ v = mn \wedge (\exists!m.\ v\$1 = m) \wedge (\exists!n.\ v\$2 =$

$n)$
using *is-int-def* **by** *auto*
let $?m = fst\ mn$
let $?n = snd\ mn$

have $?m = (THE\ m::int.\ m = v\$1)$ **using** v
by $(meson\ fstI)$
moreover have $\exists! m::int.\ m = v\$1$ **using** v *is-int-def*
by $(metis\ (no-types,\ lifting)\ mem-Collect-eq\ of-int-eq-iff)$
ultimately have m -in- $S: ?m \in S$
by $(metis\ (mono-tags,\ lifting)\ S\ mem-Collect-eq\ theI'\ v)$

have $?n = (THE\ n::int.\ n = v\$2)$ **using** v
by $(meson\ sndI)$
moreover have $\exists! n::int.\ n = v\$2$ **using** v *is-int-def*
by $(metis\ (no-types,\ lifting)\ mem-Collect-eq\ of-int-eq-iff)$

ultimately have $n\text{-in-}S$: $?n \in S$
by (*metis* (*mono-tags*, *lifting*) *S mem-Collect-eq theI' v*)

show $mn \in S \times S$ **using** $m\text{-in-}S$ $n\text{-in-}S$ v **by** *auto*
qed
ultimately show *?thesis*
by (*meson finite-S finite-SigmaI finite-imageD finite-subset*)
qed
qed
ultimately show *?thesis*
by (*smt* (*verit*) *finite-subset subsetI*)
qed

lemma *finite-path-image*:
assumes *polygon p*
shows *finite* $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$
using *bounded-finite inside-outside-polygon*
unfolding *inside-outside-def*
by (*meson assms bounded-simple-path-image polygon-def*)

lemma *finite-path-inside*:
assumes *polygon p*
shows *finite* $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$
using *bounded-finite inside-outside-polygon*
unfolding *inside-outside-def*
using *assms* **by** *presburger*

lemma *bounded-finite-inside*:
fixes $B:: (\text{real}^2)$ *set*
assumes *simple-path p*
shows *bounded* (*path-inside p*)
using *assms*
by (*simp add: bounded-inside bounded-simple-path-image path-inside-def*)

lemma *finite-integral-points-path-image*:
assumes *simple-path p*
shows *finite* $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$
using *bounded-finite bounded-simple-path-image assms* **by** *blast*

lemma *finite-integral-points-path-inside*:
assumes *simple-path p*
shows *finite* $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$
using *bounded-finite bounded-finite-inside assms* **by** *blast*

27 Pick splitting

lemma *pick-split-path-union-main*:
assumes *is-split: is-polygon-split-path vts i j cutvts*
assumes $vts1 = (\text{take } i \text{ vts})$

```

assumes vts2 = (take (j - i - 1) (drop (Suc i) vts))
assumes vts3 = drop (j - i) (drop (Suc i) vts)
assumes x = vts!i
assumes y = vts!j
assumes cutpath = make-polygonal-path (x # cutvts @ [y])
assumes p: p = make-polygonal-path (vts@[vts!0]) (is p = make-polygonal-path
?p-vts)
assumes p1: p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x]))
(is p1 = make-polygonal-path ?p1-vts)
assumes p2: p2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @
[vts ! 0]) (is p2 = make-polygonal-path ?p2-vts)
assumes I1: I1 = card {x. integral-vec x ∧ x ∈ path-inside p1}
assumes B1: B1 = card {x. integral-vec x ∧ x ∈ path-image p1}
assumes I2: I2 = card {x. integral-vec x ∧ x ∈ path-inside p2}
assumes B2: B2 = card {x. integral-vec x ∧ x ∈ path-image p2}
assumes I: I = card {x. integral-vec x ∧ x ∈ path-inside p}
assumes B: B = card {x. integral-vec x ∧ x ∈ path-image p}
assumes all-integral-vts: all-integral vts
shows measure lebesgue (path-inside p1) = I1 + B1/2 - 1
    ⇒ measure lebesgue (path-inside p2) = I2 + B2/2 - 1
    ⇒ measure lebesgue (path-inside p) = I + B/2 - 1
    measure lebesgue (path-inside p) = I + B/2 - 1
    ⇒ measure lebesgue (path-inside p2) = I2 + B2/2 - 1
    ⇒ measure lebesgue (path-inside p1) = I1 + B1/2 - 1
    measure lebesgue (path-inside p) = I + B/2 - 1
    ⇒ measure lebesgue (path-inside p1) = I1 + B1/2 - 1
    ⇒ measure lebesgue (path-inside p2) = I2 + B2/2 - 1
proof -
let ?p-im = {x. integral-vec x ∧ x ∈ path-image p}
let ?p1-im = {x. integral-vec x ∧ x ∈ path-image p1}
let ?p2-im = {x. integral-vec x ∧ x ∈ path-image p2}
let ?p-int = {x. integral-vec x ∧ x ∈ path-inside p}
let ?p1-int = {x. integral-vec x ∧ x ∈ path-inside p1}
let ?p2-int = {x. integral-vec x ∧ x ∈ path-inside p2}

have vts: vts = vts1 @ (x # (vts2 @ y # vts3))
  using assms split-up-a-list-into-3-parts
  using is-polygon-split-path-def by blast
have polygon p
  using finite-path-image assms(1) p unfolding is-polygon-split-path-def
  by (smt (verit, best))
then have B-finite: finite ?p-im
  using finite-path-image by auto
have polygon-p1: polygon p1
  using finite-path-image assms(1) p1 unfolding is-polygon-split-path-def
  by (smt (z3) assms(3) assms(5) assms(6))
then have B1-finite: finite ?p1-im
  using finite-path-image by auto
have polygon-p2: polygon p2

```

```

    using finite-path-image assms(1) p1 unfolding is-polygon-split-path-def
    by (smt (z3) assms(2) assms(4) assms(5) assms(6) p2)
then have B2-finite: finite ?p2-im
    using finite-path-image by auto

have vts-distinct: distinct vts
    using simple-polygonal-path-vts-distinct
    by (metis ‹polygon p› butlast-snoc p polygon-def)
then have x-neq-y: x ≠ y
    by (metis assms(1) assms(5) assms(6) index-first index-nth-id is-polygon-split-path-def)
then have card-2: card {x, y} = 2
    by auto
have polygon-split-props: (is-polygon-cut-path (vts@[vts!0]) cutpath ∧
    polygon p ∧ polygon p1 ∧ polygon p2 ∧
    path-inside p1 ∩ path-inside p2 = {} ∧
    path-inside p1 ∪ path-inside p2 ∪ (path-image cutpath - {x, y}) = path-inside
p
    ∧ ((path-image p1) - (path-image cutpath)) ∩ ((path-image p2) - (path-image
cutpath)) = {}
    ∧ path-image p = ((path-image p1) - (path-image cutpath)) ∪ ((path-image p2)
- (path-image cutpath)) ∪ {x, y})
    using assms
    by (meson is-polygon-split-path-def)
have measure-sum: measure lebesgue (path-inside p) = measure lebesgue (path-inside
p1) + measure lebesgue (path-inside p2)
    using polygon-split-path-add-measure assms
    by (smt (verit, del-insts))

let ?yx-int = {k. integral-vec k ∧ k ∈ path-image (make-polygonal-path (y#rev
cutvts@[x]))}
let ?xy-int = {k. integral-vec k ∧ k ∈ path-image cutpath}
have yx-int-is-xy-int: ?yx-int = ?xy-int
    using rev-vts-path-image[of x # cutvts @ [y]] assms(7) by simp
have x # vts2 @ [y] @ rev cutvts @ [x] = (x#vts2) @ ([y] @ rev cutvts @ [x]) @
[]
    by simp
then have sublist ([y]@rev cutvts@[x]) ?p1-vts
    unfolding sublist-def by blast
then have subset1:
    ?xy-int ⊆ ?p1-im
    using sublist-integral-subset-integral-on-path p1 yx-int-is-xy-int
    by force
have len-gteq: length (x # cutvts @ [y]) ≥ 2
    by auto
have sublist-p2: sublist (x # cutvts @ [y]) ?p2-vts
    unfolding sublist-def by auto
then have subset2:
    ?xy-int ⊆ ?p2-im

```

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using sublist-integral-subset-integral-on-path[OF len-gteq p2 sublist-p2]
assms(7) by blast

let ?S1 = ?p1-im - ?xy-int
let ?S2 = ?p2-im - ?xy-int
have disjoint-1: ?S1 ∩ ?S2 = {}
using polygon-split-props by blast

have integral-xy: integral-vec x ∧ integral-vec y
using all-integral-pts pts
using all-integral-def by auto
have nonempty: y # rev cutvts @ [x] ≠ []
by simp
have trivial: make-polygonal-path (y # rev cutvts @ [x]) = make-polygonal-path
(y # rev cutvts @ [x])
by auto
have pathstart (make-polygonal-path (y#rev cutvts@[x])) = y ∧ pathfinish (make-polygonal-path
(y#rev cutvts@[x])) = x
using polygon-pathstart[OF nonempty trivial] polygon-pathfinish[OF nonempty
trivial]
by (metis last.simps last-conv-nth nonempty nth-Cons-0 snoc-eq-iff-butlast)
then have x-in-y-in: x ∈ path-image (make-polygonal-path (y#rev cutvts@[x]))
∧ y ∈ path-image (make-polygonal-path (y#rev cutvts@[x]))
unfolding pathstart-def pathfinish-def path-image-def
by (metis ⟨pathstart (make-polygonal-path (y # rev cutvts @ [x])) = y ∧
pathfinish (make-polygonal-path (y # rev cutvts @ [x])) = x⟩ path-image-def pathfin-
ish-in-path-image pathstart-in-path-image)
then have {x, y} ⊆ ?yx-int
using integral-xy
by simp
then have disjoint-2: (?S1 ∪ ?S2) ∩ {x, y} = {}
by (simp add: yx-int-is-xy-int)
have path-image p =
path-image p1 - path-image cutpath ∪
(path-image p2 - path-image cutpath) ∪
{x, y}
using polygon-split-props by auto
then have set-union: ?p-im = (?S1 ∪ ?S2) ∪ {x, y}
using polygon-split-props integral-xy by auto
then have add-card: B = card (?p1-im - ?xy-int) + card (?p2-im - ?xy-int)
+ card {x, y}
using B-finite using disjoint-1 disjoint-2
by (metis (no-types, lifting) B card-Un-disjoint finite-Un)
have sub1: card (?p1-im - ?xy-int) = B1 - card ?xy-int
using B1-finite B1 subset1
by (meson card-Diff-subset finite-subset)
have sub2: card (?p2-im - ?xy-int) = B2 - card ?xy-int
using B2-finite B2 subset2
by (meson card-Diff-subset finite-subset)

```

```

have B: B = (B1 - card ?xy-int) + (B2 - card ?xy-int) + card {x, y}
  using add-card sub1 sub2
  by auto
then have B-sum-h: B = B1 + B2 - 2*card ?xy-int + 2
  using card-2
  by (smt (verit, best) B1 B1-finite B2 B2-finite Nat.add-diff-assoc add.commute
card-mono diff-diff-left mult-2 subset1 subset2)
  then have B1 + B2 = B + 2*card ?xy-int - 2
  by (metis (no-types, lifting) B1 B1-finite B2 B2-finite add-mono-thms-linordered-semiring(1)
card-mono diff-add-inverse2 le-add2 mult-2 ordered-cancel-comm-monoid-diff-class.add-diff-assoc2
subset1 subset2)
  then have B-sum: (B1 + B2)/2 = B/2 + card ?xy-int - 1
  by (smt (verit) B-sum-h field-sum-of-halves le-add2 mult-2 nat-1-add-1 of-nat-1
of-nat-add of-nat-diff ordered-cancel-comm-monoid-diff-class.add-diff-assoc2)
  have casting-h:  $\bigwedge A B:: \text{nat}. A \geq B \implies \text{real } (A - B) = \text{real } A - \text{real } B$ 
  by auto
  have path-inside p1  $\cup$  path-inside p2  $\cup$  (path-image cutpath - {x, y}) =
path-inside p
  using polygon-split-props by auto
  then have interior-union: ?p-int = (?xy-int - {x, y})  $\cup$  ?p1-int  $\cup$  ?p2-int
  by blast

have finite-inside-p: finite ?p-int
  using bounded-finite inside-outside-polygon
  by (simp add: polygon-split-props inside-outside-def)
have finite-pathimage: finite (?xy-int - {x, y})
  using B1-finite finite-subset subset1 by auto
have finite-inside-p1: finite ?p1-int
  using polygon-split-props bounded-finite inside-outside-polygon
  using finite-Un finite-inside-p interior-union by auto
have finite-inside-p2: finite ?p2-int
  using polygon-split-props bounded-finite inside-outside-polygon
  using finite-Un finite-inside-p interior-union by auto
have path-image-inside-disjoint1: (?xy-int - {x, y})  $\cap$  (?p1-int) = {}
  using subset1 inside-outside-polygon[OF polygon-p1]
  unfolding inside-outside-def by auto
have path-image-inside-disjoint2: (?xy-int - {x, y})  $\cap$  (?p2-int) = {}
  using subset2 inside-outside-polygon[OF polygon-p2]
  unfolding inside-outside-def by auto

have (?xy-int - {x, y})  $\cap$  (?p1-int  $\cup$  ?p2-int) = {}
  using subset2 path-image-inside-disjoint1 path-image-inside-disjoint2
  by auto
then have I-is: I = card (?xy-int - {x, y}) +
card (?p1-int  $\cup$  ?p2-int)
  using interior-union I finite-inside-p1 finite-inside-p2
  by (metis (no-types, lifting) card-Un-disjoint finite-Un finite-pathimage sup-assoc)

have disjoint-4: ?p1-int  $\cap$  ?p2-int = {}

```

```

    using polygon-split-props by auto
  then have I = card (?xy-int - {x, y}) +
    I1 + I2
    using I-is finite-inside-p1 finite-inside-p2
    by (simp add: I1 I2 card-Un-disjoint)
  have interior-subset: (?xy-int - {x, y})  $\subseteq$  ?p-int
    using interior-union by auto
  have x-y-subset: {x, y}  $\subseteq$  ?xy-int
    using x-in-y-in rev-vts-path-image[of x # cutvts @ [y]] assms(7)
    integral-xy
    using yx-int-is-xy-int by blast
  have real (card (?xy-int - {x, y})) =
    real (card (?xy-int)) - real (card {x, y})
    using x-y-subset
    by (metis (no-types, lifting) B2-finite card-Diff-subset card-mono finite-subset
of-nat-diff subset2)
  then have card-diff: real (card (?xy-int - {x, y})) =
    real (card (?xy-int)) - 2
    using card-2 by auto
  then have I = I1 + I2 + (card (?xy-int - {x, y}))
    using I I1 I2 interior-union finite-inside-p1 finite-inside-p2
    by (simp add: I-is disjoint-4 card-Un-disjoint)
  then have I = I1 + I2 + real (card (?xy-int)) - 2
    using card-diff
    by linarith
  then have I-sum: I1 + I2 = I - real (card ?xy-int) + 2
    by fastforce

{assume pick1: measure lebesgue (path-inside p1) = I1 + B1/2 - 1
  assume pick2: measure lebesgue (path-inside p2) = I2 + B2/2 - 1
  have measure lebesgue (path-inside p) = I1 + I2 + (B1+B2)/2 - 2
    using pick1 pick2 measure-sum by auto
  then have measure lebesgue (path-inside p) = I - real (card ?xy-int) + 2 +
    B/2 + card ?xy-int - 1 - 2
    using I-sum B-sum
    by linarith
  then have measure lebesgue (path-inside p) = I + B/2 - 1 by auto
}
  then show measure lebesgue (path-inside p1) = I1 + B1/2 - 1  $\implies$  measure
lebesgue (path-inside p2) = I2 + B2/2 - 1  $\implies$  measure lebesgue (path-inside p)
= I + B/2 - 1
    by blast

{assume pick1: measure lebesgue (path-inside p) = I + B/2 - 1
  assume pick2: measure lebesgue (path-inside p2) = I2 + B2/2 - 1
  then have real I + real B / 2 - 1 = (measure lebesgue (path-inside p1)) +
I2 + B2/2 - 1
    using measure-sum pick1 pick2 by auto
  then have measure lebesgue (path-inside p) = I - real (card ?xy-int) + 2 +

```

```

    B/2 + card ?xy-int - 1 - 2
    using I-sum B-sum pick1
    by linarith
  then have measure lebesgue (path-inside p1) = I1 + B1/2 - 1
    using B-sum ⟨real I = real (I1 + I2) + real (card {k. integral-vec k ∧ k ∈
path-image cutpath}) - 2⟩ field-sum-of-halves measure-sum of-nat-add
    pick1 pick2 by auto
  }
  then show measure lebesgue (path-inside p) = I + B/2 - 1 ⇒ measure
lebesgue (path-inside p2) = I2 + B2/2 - 1 ⇒ measure lebesgue (path-inside p1)
= I1 + B1/2 - 1
    by blast

{assume pick1: measure lebesgue (path-inside p) = I + B/2 - 1
  assume pick2: measure lebesgue (path-inside p1) = I1 + B1/2 - 1
  then have real I + real B / 2 - 1 = (measure lebesgue (path-inside p2)) +
I1 + B1/2 - 1
    using measure-sum pick1 pick2 by auto
  then have measure lebesgue (path-inside p) = I - real (card ?xy-int) + 2 +
B/2 + card ?xy-int - 1 - 2
    using I-sum B-sum pick1
    by linarith
  then have measure lebesgue (path-inside p2) = I2 + B2/2 - 1
    using B-sum ⟨real I = real (I1 + I2) + real (card {k. integral-vec k ∧ k ∈
path-image cutpath}) - 2⟩ field-sum-of-halves measure-sum of-nat-add
    using pick2 by auto
  }
  then show measure lebesgue (path-inside p) = I + B/2 - 1 ⇒ measure lebesgue
(path-inside p1) = I1 + B1/2 - 1 ⇒ measure lebesgue (path-inside p2) = I2 +
B2/2 - 1
    by blast
qed

```

lemma *pick-split-union*:

```

  assumes is-split: is-polygon-split vts i j
  assumes vts1 = (take i vts)
  assumes vts2 = (take (j - i - 1) (drop (Suc i) vts))
  assumes vts3 = drop (j - i) (drop (Suc i) vts)
  assumes x = vts ! i
  assumes y = vts ! j
  assumes p: p = make-polygonal-path (vts@[vts!0]) (is p = make-polygonal-path
?p-vts)
  assumes p1: p1 = make-polygonal-path (x#(vts2@[y, x])) (is p1 = make-polygonal-path
?p1-vts)
  assumes p2: p2 = make-polygonal-path (vts1 @ [x, y] @ vts3 @ [vts ! 0]) (is p2
= make-polygonal-path ?p2-vts)
  assumes I1: I1 = card {x. integral-vec x ∧ x ∈ path-inside p1}
  assumes B1: B1 = card {x. integral-vec x ∧ x ∈ path-image p1}
  assumes pick1: measure lebesgue (path-inside p1) = I1 + B1/2 - 1

```

```

assumes I2: I2 = card {x. integral-vec x ∧ x ∈ path-inside p2}
assumes B2: B2 = card {x. integral-vec x ∧ x ∈ path-image p2}
assumes pick2: measure lebesgue (path-inside p2) = I2 + B2/2 - 1
assumes I: I = card {x. integral-vec x ∧ x ∈ path-inside p}
assumes B: B = card {x. integral-vec x ∧ x ∈ path-image p}
assumes all-integral-vts: all-integral vts
shows measure lebesgue (path-inside p) = I + B/2 - 1
        measure lebesgue (path-inside p) = measure lebesgue (path-inside p1) +
measure lebesgue (path-inside p2)
proof -
let ?p-im = {x. integral-vec x ∧ x ∈ path-image p}
let ?p1-im = {x. integral-vec x ∧ x ∈ path-image p1}
let ?p2-im = {x. integral-vec x ∧ x ∈ path-image p2}
let ?p-int = {x. integral-vec x ∧ x ∈ path-inside p}
let ?p1-int = {x. integral-vec x ∧ x ∈ path-inside p1}
let ?p2-int = {x. integral-vec x ∧ x ∈ path-inside p2}

have vts: vts = vts1 @ (x # (vts2 @ y # vts3))
  using assms split-up-a-list-into-3-parts
  using is-polygon-split-def by blast
have polygon p
  using finite-path-image assms(1) p unfolding is-polygon-split-def
  by (smt (verit, best))
then have B-finite: finite ?p-im
  using finite-path-image by auto
have polygon-p1: polygon p1
  using finite-path-image assms(1) p1 unfolding is-polygon-split-def
  by (smt (z3) assms(3) assms(5) assms(6))
then have B1-finite: finite ?p1-im
  using finite-path-image by auto
have polygon-p2: polygon p2
  using finite-path-image assms(1) p1 unfolding is-polygon-split-def
  by (smt (z3) assms(2) assms(4) assms(5) assms(6) p2)
then have B2-finite: finite ?p2-im
  using finite-path-image by auto

have vts-distinct: distinct vts
  using simple-polygonal-path-vts-distinct
  by (metis ⟨polygon p⟩ butlast-snoc p polygon-def)
then have x-neq-y: x ≠ y
  by (metis assms(1) assms(5) assms(6) index-first index-nth-id is-polygon-split-def)
then have card-2: card {x, y} = 2
  by auto
have polygon-split-props: is-polygon-cut ?p-vts x y ∧
  polygon p ∧ polygon p1 ∧ polygon p2 ∧
  path-inside p1 ∩ path-inside p2 = {} ∧
  path-inside p1 ∪ path-inside p2 ∪ (path-image (linepath x y) - {x, y})
  = path-inside p ∧ ((path-image p1) - (path-image (linepath x y))) ∩
  ((path-image p2) - (path-image (linepath x y))) = {}

```

$\wedge \text{path-image } p = ((\text{path-image } p1) - (\text{path-image } (\text{linepath } x \ y))) \cup ((\text{path-image } p2) - (\text{path-image } (\text{linepath } x \ y))) \cup \{x, y\}$
using *assms*
by (*meson is-polygon-split-def*)
have *measure lebesgue (path-inside p) = measure lebesgue (path-inside p1) + measure lebesgue (path-inside p2)*
using *polygon-split-add-measure assms*
by (*smt (verit, del-insts)*)
then have *measure-sum: measure lebesgue (path-inside p) = I1 + I2 + (B1+B2)/2*
 -2
using *pick1 pick2 by auto*

let *?yx-int = {k. integral-vec k \wedge k \in path-image (linepath y x)}*
let *?xy-int = {k. integral-vec k \wedge k \in path-image (linepath x y)}*
have *yx-int-is-xy-int: ?yx-int = ?xy-int*
by (*simp add: closed-segment-commute*)

have *sublist [y, x] ?p1-vts by (simp add: sublist-Cons-right)*
then have *subset1:*
?xy-int \subseteq ?p1-im
using *sublist-pair-integral-subset-integral-on-path p1 yx-int-is-xy-int by blast*
have *subset2:*
?xy-int \subseteq ?p2-im
using *sublist-pair-integral-subset-integral-on-path p2 by blast*

let *?S1 = ?p1-im - ?xy-int*
let *?S2 = ?p2-im - ?xy-int*
have *disjoint-1: ?S1 \cap ?S2 = {}*
using *polygon-split-props by blast*

have *integral-xy: integral-vec x \wedge integral-vec y*
using *all-integral-vts vts*
using *all-integral-def by auto*
then have *{x, y} \subseteq ?yx-int*
by *simp*
then have *disjoint-2: (?S1 \cup ?S2) \cap {x, y} = {}*
by *simp*
have *path-image p =*
path-image p1 - path-image (linepath x y) \cup
(path-image p2 - path-image (linepath x y)) \cup
{x, y}
using *polygon-split-props by auto*
then have *set-union: ?p-im = (?S1 \cup ?S2) \cup {x, y}*
using *polygon-split-props integral-xy by auto*
then have *add-card: B = card (?p1-im - ?xy-int) + card (?p2-im - ?xy-int)*
 $+ \text{card } \{x, y\}$
using *B-finite using disjoint-1 disjoint-2*
by (*metis (no-types, lifting) B card-Un-disjoint finite-Un*)
have *sub1: card (?p1-im - ?xy-int) = B1 - card ?xy-int*

```

using B1-finite B1 subset1
by (meson card-Diff-subset finite-subset)
have sub2: card (?p2-int - ?xy-int) = B2 - card ?xy-int
using B2-finite B2 subset2
by (meson card-Diff-subset finite-subset)
have B: B = (B1 - card ?xy-int) + (B2 - card ?xy-int) + card {x, y}
using add-card sub1 sub2
by auto
then have B-sum-h: B = B1 + B2 - 2*card ?xy-int + 2
using card-2
by (smt (verit, best) B1 B1-finite B2 B2-finite Nat.add-diff-assoc add.commute
card-mono diff-diff-left mult-2 subset1 subset2)
then have B1 + B2 = B + 2*card ?xy-int - 2
by (metis (no-types, lifting) B1 B1-finite B2 B2-finite add-mono-thms-linordered-semiring(1)
card-mono diff-add-inverse2 le-add2 mult-2 ordered-cancel-comm-monoid-diff-class.add-diff-assoc2
subset1 subset2)
then have B-sum: (B1 + B2)/2 = B/2 + card ?xy-int - 1
by (smt (verit) B-sum-h field-sum-of-halves le-add2 mult-2 nat-1-add-1 of-nat-1
of-nat-add of-nat-diff ordered-cancel-comm-monoid-diff-class.add-diff-assoc2)
have casting-h:  $\bigwedge A B:: \text{nat}. A \geq B \implies \text{real } (A - B) = \text{real } A - \text{real } B$ 
by auto
have path-inside p1  $\cup$  path-inside p2  $\cup$  (path-image (linepath x y) - {x, y}) =
path-inside p
using polygon-split-props by auto
then have interior-union: ?p-int = (?xy-int - {x, y})  $\cup$  ?p1-int  $\cup$  ?p2-int
by blast

have finite-inside-p: finite ?p-int
using bounded-finite inside-outside-polygon
by (simp add: polygon-split-props inside-outside-def)
have finite-pathimage: finite (?xy-int - {x, y})
using B1-finite finite-subset subset1 by auto
have finite-inside-p1: finite ?p1-int
using polygon-split-props bounded-finite inside-outside-polygon
using finite-Un finite-inside-p interior-union by auto
have finite-inside-p2: finite ?p2-int
using polygon-split-props bounded-finite inside-outside-polygon
using finite-Un finite-inside-p interior-union by auto
have path-image-inside-disjoint1: (?xy-int - {x, y})  $\cap$  (?p1-int) = {}
using subset1 inside-outside-polygon[OF polygon-p1]
unfolding inside-outside-def by auto
have path-image-inside-disjoint2: (?xy-int - {x, y})  $\cap$  (?p2-int) = {}
using subset2 inside-outside-polygon[OF polygon-p2]
unfolding inside-outside-def by auto
have (?xy-int - {x, y})  $\cap$  (?p1-int  $\cup$  ?p2-int) = {}
using subset2 path-image-inside-disjoint1 path-image-inside-disjoint2
by auto
then have I-is: I = card (?xy-int - {x, y}) +
card (?p1-int  $\cup$  ?p2-int)

```

```

using interior-union I finite-inside-p1 finite-inside-p2
by (metis (no-types, lifting) card-Un-disjoint finite-Un finite-pathimage sup-assoc)

have disjoint-4: ?p1-int  $\cap$  ?p2-int = {}
using polygon-split-props by auto
then have I = card (?xy-int - {x, y}) +
  I1 + I2
using I-is finite-inside-p1 finite-inside-p2
by (simp add: I1 I2 card-Un-disjoint)
have interior-subset: (?xy-int - {x, y})  $\subseteq$  ?p-int
using interior-union by auto
have x-y-subset: {x, y}  $\subseteq$  ?xy-int
using local.set-union by auto
have real (card (?xy-int - {x, y})) =
  real (card (?xy-int)) - real (card {x, y})
using x-y-subset
by (metis (no-types, lifting) B2-finite card-Diff-subset card-mono finite-subset
of-nat-diff subset2)
then have card-diff: real (card (?xy-int - {x, y})) =
  real (card (?xy-int)) - 2
using card-2 by auto
then have I = I1 + I2 + (card (?xy-int - {x, y}))
using I I1 I2 interior-union finite-inside-p1 finite-inside-p2
by (simp add: I-is disjoint-4 card-Un-disjoint)
then have I = I1 + I2 + real (card (?xy-int)) - 2
using card-diff
by linarith
then have I-sum: I1 + I2 = I - real (card ?xy-int) + 2
by fastforce
have measure lebesgue (path-inside p) = I - real (card ?xy-int) + 2 +
  B/2 + card ?xy-int - 1 - 2
using measure-sum I-sum B-sum
by linarith
then show measure lebesgue (path-inside p) = I + B/2 - 1 by auto

show measure lebesgue (path-inside p) = measure lebesgue (path-inside p1) +
  measure lebesgue (path-inside p2)
using ‹Sigma-Algebra.measure lebesgue (path-inside p) = Sigma-Algebra.measure
  lebesgue (path-inside p1) + Sigma-Algebra.measure lebesgue (path-inside p2)› by
  blast
qed

```

```

lemma pick-split-path-union:
assumes is-split: is-polygon-split-path vts i j cutvts
assumes vts1 = (take i vts)
assumes vts2 = (take (j - i - 1) (drop (Suc i) vts))
assumes vts3 = drop (j - i) (drop (Suc i) vts)
assumes x = vts!i
assumes y = vts!j

```

```

assumes cutpath = make-polygonal-path (x # cutvts @ [y])
assumes p: p = make-polygonal-path (vts@[vts!0]) (is p = make-polygonal-path
?p-vts)
assumes p1: p1 = make-polygonal-path (x#(vts2 @ [y] @ (rev cutvts) @ [x]))
(is p1 = make-polygonal-path ?p1-vts)
assumes p2: p2 = make-polygonal-path (vts1 @ ([x] @ cutvts @ [y]) @ vts3 @
[vts ! 0]) (is p2 = make-polygonal-path ?p2-vts)
assumes I1: I1 = card {x. integral-vec x ∧ x ∈ path-inside p1}
assumes B1: B1 = card {x. integral-vec x ∧ x ∈ path-image p1}
assumes pick1: measure lebesgue (path-inside p1) = I1 + B1/2 - 1
assumes I2: I2 = card {x. integral-vec x ∧ x ∈ path-inside p2}
assumes B2: B2 = card {x. integral-vec x ∧ x ∈ path-image p2}
assumes pick2: measure lebesgue (path-inside p2) = I2 + B2/2 - 1
assumes I: I = card {x. integral-vec x ∧ x ∈ path-inside p}
assumes B: B = card {x. integral-vec x ∧ x ∈ path-image p}
assumes all-integral-vts: all-integral vts
shows measure lebesgue (path-inside p) = I + B/2 - 1
using pick-split-path-union-main pick1 pick2(1) assms by blast

```

lemma *pick-triangle-basic-split*:

```

assumes p = make-triangle a b c and distinct [a, b, c] and ¬ collinear {a, b,
c} and

```

```

d-prop: d ∈ path-image (linepath a b) ∧ d ∉ {a, b, c}

```

```

shows good-linepath c d [a, d, b, c, a]

```

```

∧ path-image (make-polygonal-path [a, d, b, c, a]) = path-image p

```

proof –

```

let ?l = linepath c d

```

```

let ?L = path-image ?l

```

```

let ?P = path-image p

```

```

let ?vts' = [a, d, b, c, a]

```

```

let ?p' = make-polygonal-path ?vts'

```

```

let ?P' = path-image ?p'

```

```

have h1: path-image (make-polygonal-path [a, b, c, a]) = path-image (linepath a
b) ∪ path-image (linepath b c) ∪ path-image (linepath c a)

```

```

using polygonal-path-image-linepath-union by (simp add: path-image-join sup.assoc)

```

```

have h2: path-image (make-polygonal-path [a, d, b, c, a]) = path-image (linepath a
d) ∪ path-image (linepath d b) ∪ path-image (linepath b c) ∪ path-image (linepath
c a)

```

```

using polygonal-path-image-linepath-union by (simp add: path-image-join sup.assoc)

```

```

have h3: path-image (linepath a b) = path-image (linepath a d) ∪ path-image
(linepath d b)

```

```

using path-image-linepath-union d-prop by auto

```

```

have 1: ?P' = ?P

```

```

using h1 h2 h3

```

```

using assms(1) make-triangle-def by force

```

```

have {c, d} = ?L ∩ ?P

```

```

proof(rule ccontr)
  have subs: {c, d} ⊆ ?L ∩ ?P
    using assms(1) vertices-on-path-image unfolding make-triangle-def
    by (metis IntD2 IntI assms(4) empty-subsetI inf-sup-absorb insert-subset
list.discI list.simps(15) nth-Cons-0 path-image-cons-union pathfinish-in-path-image
pathfinish-linepath pathstart-in-path-image pathstart-linepath)

  assume *: {c, d} ≠ ?L ∩ ?P
  then obtain z where z: z ≠ c ∧ z ≠ d ∧ z ∈ ?L ∩ ?P using subs by blast
  then have cases:
    z ∈ path-image (linepath a b) ∨ z ∈ path-image (linepath b c) ∨ z ∈ path-image
(linepath c a)
    using 1 h2 h3 by blast
    { assume **: z ∈ path-image (linepath a b)
      moreover have z ∈ ?L ∧ d ∈ ?L ∧ d ∈ path-image (linepath a b) using
assms z by force
      ultimately have {z, d} ⊆ ?L ∩ path-image (linepath a b) ∧ z ≠ d using z
by blast
      then have collinear {a, b, c, d} using two-linepath-colinearity-property by
fastforce
      then have False using assms(2) assms(3) collinear-4-3 by auto
    } moreover
    { assume **: z ∈ path-image (linepath b c)
      then have collinear {a, b, c, d} using two-linepath-colinearity-property[of z
- b c c d]
      by (smt (verit) ** IntE assms(3) collinear-3-trans d-prop in-path-image-imp-collinear
insertCI insert-commute z)
      then have False using assms(2) assms(3) collinear-4-3 by auto
    } moreover
    { assume **: z ∈ path-image (linepath c a)
      then have collinear {a, b, c, d} using two-linepath-colinearity-property[of z
- c a c d]
      by (smt (verit) IntD1 assms(3) collinear-3-trans d-prop in-path-image-imp-collinear
insert-commute insert-iff z)
      then have False using assms(2) assms(3) collinear-4-3 by auto
    }
    ultimately show False using cases by argo
  qed
  moreover have ?L ⊆ path-inside p ∪ ?P
  proof–
    have convex hull {a, b, c} = path-inside p ∪ ?P
      by (simp add: Un-commute assms(1) assms(3) triangle-convex-hull)
    moreover have ?L ⊆ convex hull {a, b, c}
      by (smt (verit, ccfv-threshold) assms empty-subsetI hull-insert hull-mono in-
sert-commute insert-mono insert-subset path-image-linepath segment-convex-hull)
    ultimately show ?thesis by blast
  qed
  ultimately have ?L ⊆ path-inside p ∪ {c, d} by blast
  then have ?L ⊆ path-inside ?p' ∪ {c, d} using 1 unfolding path-inside-def by

```

presburger

then have 2: *good-linepath* $c\ d\ ?vts'$ **using** *assms unfolding good-linepath-def*
by *auto*

thus *?thesis* **using** 1 **by** *blast*
qed

28 Convex Hull Has Good Linepath

lemma *leq-2-extreme-points-means-collinear*:

fixes $vts :: 'a::euclidean-space\ set$

assumes *finite vts*

assumes $\text{card } \{v. v\ \text{extreme-point-of } (\text{convex hull } vts)\} \leq 2$

shows *collinear vts*

using *assms*

by (*metis Krein-Milman-polytope affine-hull-convex-hull collinear-affine-hull-collinear collinear-small extreme-points-of-convex-hull finite-subset*)

lemma *convex-hull-non-extreme-point-in-open-seg*:

assumes $H = \text{convex hull } vts$

assumes $x \in H - \{v. v\ \text{extreme-point-of } H\}$

shows $\exists a\ b. a \in H \wedge b \in H \wedge x \in \text{open-segment } a\ b$

using *assms unfolding extreme-point-of-def* **by** *blast*

lemma *convex-hull-extreme-points-vertex-split*:

fixes $vts :: (\mathbb{R}^2)\ set$

assumes $H = \text{convex hull } vts$

assumes *finite vts*

assumes $\text{card } \{v. v\ \text{extreme-point-of } H\} \geq 4$

assumes $\{a, b, c\} \subseteq \{v. v\ \text{extreme-point-of } H\} \wedge \text{distinct } [a, b, c]$

shows $\text{path-image } (\text{linepath } a\ b) \cap \text{interior } H \neq \{\}$

$\vee \text{path-image } (\text{linepath } b\ c) \cap \text{interior } H \neq \{\}$

$\vee \text{path-image } (\text{linepath } c\ a) \cap \text{interior } H \neq \{\}$

proof –

let $?ep = \{v. v\ \text{extreme-point-of } H\}$

have $H: H = \text{convex hull } ?ep$ **using** *Krein-Milman-polytope assms(1) assms(2)*

by *blast*

let $?H' = \text{convex hull } \{a, b, c\}$

have *not-collinear*: $\neg \text{collinear } \{a, b, c\}$

proof(*rule ccontr*)

assume $\neg \neg \text{collinear } \{a, b, c\}$

then have *collinear* $\{a, b, c\}$ **by** *blast*

then have $a \in \text{path-image } (\text{linepath } b\ c)$

$\vee b \in \text{path-image } (\text{linepath } a\ c)$

$\vee c \in \text{path-image } (\text{linepath } a\ b)$

using *collinear-between-cases unfolding between-def*

by (*smt (verit, del-insts) between-mem-segment closed-segment-eq collinear-between-cases*)

doubleton-eq-iff path-image-linepath)
moreover have $a \neq b \wedge b \neq c \wedge a \neq c$ **using** *assms* **by** *simp*
ultimately have $a \in \text{open-segment } b \ c \vee b \in \text{open-segment } a \ c \vee c \in$
open-segment } a \ b
using *closed-segment-eq-open* **by** *auto*
moreover have $a \text{ extreme-point-of } H \wedge b \text{ extreme-point-of } H \wedge c \text{ extreme-point-of}$
H
using *assms* **by** *blast*
ultimately show *False* **unfolding** *extreme-point-of-def* **by** *blast*
qed

have *strict-subset: interior ?H' \subset interior H*
proof–
have *interior ?H' \subseteq interior H*
by (*metis H assms(4) hull-mono interior-mono*)
moreover have $?H' \subset H$
proof–
have $\text{card } \{a, b, c\} \leq 3$
by (*metis card.empty card-insert-disjoint collinear-2 finite.emptyI finite-insert*
insert-absorb nat-le-linear not-collinear numeral-3-eq-3)
then have $\text{card } (?ep - \{a, b, c\}) \geq 1$
using *assms(3) assms(4)* **by** *auto*
then obtain d **where** $d \in ?ep - \{a, b, c\}$
by (*metis One-nat-def all-not-in-conv card.empty not-less-eq-eq zero-le*)
thus *?thesis*
by (*metis DiffE H assms(4) extreme-point-of-convex-hull hull-mono mem-Collect-eq*
order-less-le)
qed
ultimately show *?thesis*
by (*metis (no-types, lifting) assms(1) assms(2) closure-convex-hull con-*
conv-closure-rel-interior convex-convex-hull convex-hull-eq-empty convex-polygon-frontier-is-path-image2
dual-order.strict-iff-order finite.emptyI finite.insertI finite.imp-bounded-convex-hull
finite.imp-compact frontier-empty insert-not-empty inside-frontier-eq-interior not-collinear
path-inside-def polygon-frontier-is-path-image rel-interior-nonempty-interior sup-bot.right-neutral
triangle-convex-hull triangle-is-convex triangle-is-polygon)
qed
moreover have *interior ?H' \neq {}*
by (*metis not-collinear convex-convex-hull convex-hull-eq-empty convex-polygon-frontier-is-path-image2*
finite.emptyI finite.insertI finite.imp-bounded-convex-hull frontier-empty insert-not-empty
inside-frontier-eq-interior path-inside-def polygon-frontier-is-path-image sup-bot.right-neutral
triangle-convex-hull triangle-is-convex triangle-is-polygon)
ultimately obtain $x \ y$ **where** $xy: x \in \text{interior } ?H' \wedge y \in \text{interior } H - \text{interior}$
 $?H'$ **by** *blast*

let $?l = \text{linepath } x \ y$

have $x \in \text{interior } ?H' \wedge y \in -(\text{interior } ?H')$ **using** xy **by** *blast*
then have $\text{path-image } ?l \cap \text{interior } ?H' \neq \{\} \wedge \text{path-image } ?l \cap -(\text{interior } ?H')$
 $\neq \{\}$ **by** *auto*

moreover have *path-connected* (*interior* ?*H'*) **by** (*simp add: convex-imp-path-connected*)
ultimately obtain *z* **where** *z*: *z* ∈ *path-image* ?*l* ∩ *frontier* (*interior* ?*H'*)
by (*metis Diff-eq Diff-eq-empty-iff all-not-in-conv convex-convex-hull convex-imp-path-connected*
path-connected-not-frontier-subset path-image-linepath segment-convex-hull)
moreover have *path-image* ?*l* ⊆ *interior* *H* **using** *xy convex-interior*[*of H*]
by (*metis DiffD1 IntD2 strict-subset assms(1) closed-segment-subset convex-convex-hull*
inf.strict-order-iff path-image-linepath)
ultimately have *z-interior*: *z* ∈ *interior* *H* **by** *blast*

have *z* ∈ *frontier* (*interior* ?*H'*) **using** *z* **by** *blast*
moreover have *frontier* (*interior* ?*H'*)
= *path-image* (*linepath* *a* *b*) ∪ *path-image* (*linepath* *b* *c*) ∪ *path-image* (*linepath*
c *a*)
proof–
let ?*p* = *make-triangle* *a* *b* *c*
have *path-inside* ?*p* = *interior* ?*H'*
by (*metis not-collinear bounded-convex-hull bounded-empty bounded-insert con-*
vox-convex-hull convex-polygon-frontier-is-path-image2 inside-frontier-eq-interior path-inside-def
triangle-convex-hull triangle-is-convex triangle-is-polygon)
then have *path-image* ?*p* = *frontier* (*interior* ?*H'*)
by (*metis not-collinear polygon-frontier-is-path-image triangle-is-polygon*)
moreover have *path-image* ?*p*
= *path-image* (*linepath* *a* *b*) ∪ *path-image* (*linepath* *b* *c*) ∪ *path-image* (*linepath*
c *a*)
by (*metis Un-assoc list.discI make-polygonal-path.simps(3) make-triangle-def*
nth-Cons-0 path-image-cons-union)
ultimately show ?*thesis* **by** *presburger*
qed
ultimately show ?*thesis* **using** *z-interior* **by** *blast*
qed

lemma *convex-hull-has-vertex-split-helper-wlog*:

assumes *p* = *make-triangle* *a* *b* *c* **and** *distinct* [*a*, *b*, *c*] **and** ¬ *collinear* {*a*, *b*,
c} **and**

d-prop: *d* ∈ *path-image* (*linepath* *a* *b*) ∧ *d* ∉ {*a*, *b*, *c*}

shows *path-image* (*linepath* *c* *d*) ∩ *path-inside* *p* ≠ {}

proof–

have *good-linepath* *c* *d* [*a*, *d*, *b*, *c*, *a*]

∧ *path-image* (*make-polygonal-path* [*a*, *d*, *b*, *c*, *a*]) = *path-image* *p*

using *pick-triangle-basic-split*[*of p a b c d*] *assms* **by** *fast*

thus ?*thesis*

unfolding *good-linepath-def*

by (*smt (verit, del-insts) Int-Un-eq(4) Int-insert-right-if1 Un-insert-right diff-points-path-image-set-property*
le-iff-inf path-inside-def pathfinish-in-path-image pathfinish-linepath pathstart-in-path-image
pathstart-linepath)

qed

lemma *convex-hull-has-vertex-split-helper*:

assumes *p* = *make-triangle* *a* *b* *c* **and** *distinct* [*a*, *b*, *c*] **and** ¬ *collinear* {*a*, *b*,

```

c} and
  d-prop:  $d \in \text{path-image } p \wedge d \notin \{a, b, c\}$ 
  shows  $\exists x y. \{x, y\} \subseteq \{a, b, c, d\} \wedge x \neq y \wedge \text{path-image } (\text{linepath } x \ y) \cap \text{path-inside } p \neq \{\}$ 
proof -
  { assume  $d \in \text{path-image } (\text{linepath } a \ b)$ 
    then have ?thesis
      using convex-hull-has-vertex-split-helper-wlog[of p a b c d] assms(1) assms(2)
    assms(3) d-prop
      by fastforce
  } moreover
  { assume *:  $d \in \text{path-image } (\text{linepath } b \ c)$ 
    let ?p' = make-triangle b c a
    have  $\text{path-image } (\text{linepath } a \ d) \cap \text{path-inside } ?p' \neq \{\}$ 
      using convex-hull-has-vertex-split-helper-wlog[of ?p' b c a d]
      by (metis (no-types, opaque-lifting) * assms(3) collinear-2 d-prop distinct-length-2-or-more
distinct-singleton insert-absorb2 insert-commute)
    moreover have  $\text{path-inside } ?p' = \text{path-inside } p$ 
      unfolding make-triangle-def
      by (smt (verit, best) assms(1) assms(3) convex-polygon-frontier-is-path-image2
insert-commute make-triangle-def path-inside-def triangle-convex-hull triangle-is-convex
triangle-is-polygon)
    ultimately have ?thesis using assms by auto
  } moreover
  { assume *:  $d \in \text{path-image } (\text{linepath } c \ a)$ 
    let ?p' = make-triangle c a b
    have  $\text{path-image } (\text{linepath } b \ d) \cap \text{path-inside } ?p' \neq \{\}$ 
      using convex-hull-has-vertex-split-helper-wlog[of ?p' c a b d]
      by (metis (no-types, opaque-lifting) * assms(3) collinear-2 d-prop distinct-length-2-or-more
distinct-singleton insert-absorb2 insert-commute)
    moreover have  $\text{path-inside } ?p' = \text{path-inside } p$ 
      unfolding make-triangle-def
      by (smt (verit, ccfv-SIG) assms(1) assms(3) convex-polygon-frontier-is-path-image2
insert-commute make-triangle-def path-inside-def triangle-convex-hull triangle-is-convex
triangle-is-polygon)
    ultimately have ?thesis using assms by auto
  }
}
ultimately show ?thesis using on-triangle-path-image-cases assms(1) d-prop
by fast
qed

```

```

lemma convex-hull-has-vertex-split:
  fixes vts :: (real^2) set
  assumes H = convex hull vts
  assumes  $\neg \text{collinear } vts$ 
  assumes card vts > 3
  assumes finite vts
  shows  $\exists a b. \{a, b\} \subseteq vts \wedge a \neq b \wedge \text{path-image } (\text{linepath } a \ b) \cap \text{interior } H \neq \{\}$ 

```

proof–
let $?ep = \{v. v \text{ extreme-point-of } H\}$
have $ep: ?ep \subseteq vts$ **by** (*simp add: assms(1) extreme-points-of-convex-hull*)
have $card-ep: card ?ep \geq 3$
by (*metis One-nat-def Suc-1 assms(1) assms(2) assms(3) card.infinite leq-2-extreme-points-means-collinear not-less-eq-eq not-less-zero numeral-3-eq-3*)
obtain $a b c$ **where** $abc: \{a, b, c\} \subseteq ?ep \wedge a \neq b \wedge b \neq c \wedge a \neq c$
proof–
obtain $a A$ **where** $a \in ?ep \wedge A = ?ep - \{a\} \wedge card A \geq 2$ **using** $card-ep$ **by**
force
moreover then obtain $b B$ **where** $b \in A \wedge B = A - \{b\} \wedge card B \geq 1$
by (*metis Suc-1 Suc-diff-le bot.extremum-uniqueI bot-nat-0.extremum card-Diff-singleton card-eq-0-iff diff-Suc-1 less-Suc-eq-le less-one linorder-not-le subset-emptyI*)
moreover then obtain $c C$ **where** $c \in B \wedge C = B - \{c\} \wedge card C \geq 0$
by (*metis One-nat-def bot-nat-0.extremum card.empty equals0I not-less-eq-eq*)
ultimately have $\{a, b, c\} \subseteq ?ep \wedge a \neq b \wedge b \neq c \wedge a \neq c$ **by** *blast*
thus *?thesis using that by auto*
qed
{ assume $*$: $card ?ep = 3$
then have $abc: ?ep = \{a, b, c\}$
by (*metis abc card-3-iff card-gt-0-iff numeral-3-eq-3 order-less-le psubset-card-mono zero-less-Suc*)
obtain d **where** $d: d \in vts \wedge d \neq a \wedge d \neq b \wedge d \neq c$
by (*metis * assms(3) abc ep insertCI nat-less-le subsetI subset-antisym*)
{ assume $d \in interior H$
then have $d \in path-image (linepath a d) \cap interior H$ **by** *simp*
then have *?thesis using ep abc d by auto*
} **moreover**
{ assume $***: d \notin interior H$
let $?p = make-triangle a b c$
have $H: H = convex hull ?ep$
proof–
have *compact H*
by (*metis assms(1) assms(3) card-eq-0-iff finite-imp-compact-convex-hull gr-implies-not0*)
moreover have *convex H using convex-convex-hull[of vts] assms by blast*
ultimately have $H = closure (convex hull ?ep)$ **using** *Krein-Milman[of H]*
by *fast*
thus *?thesis using abc by auto*
qed
then have $interior: path-inside ?p = interior H$
using abc
by (*metis assms(1,2) affine-hull-convex-hull collinear-affine-hull-collinear convex-convex-hull convex-polygon-frontier-is-path-image2 finite.intros(1) finite-imp-bounded-convex-hull finite-insert inside-frontier-eq-interior path-inside-def triangle-convex-hull triangle-is-convex triangle-is-polygon*)
then have $d-frontier: d \in frontier H$
by (*metis *** Diff-iff assms(1) UnCI d closure-Un-frontier frontier-def hull-subset in-mono*)

moreover have $\text{path-image } ?p = \text{frontier } H$
using *convex-polygon-frontier-is-path-image*
by (*metis* *assms*(1,2) *H abc affine-hull-convex-hull collinear-affine-hull-collinear convex-polygon-frontier-is-path-image2 triangle-convex-hull triangle-is-convex triangle-is-polygon*)
ultimately have $d \in \text{path-image } ?p$ **by** *blast*
moreover have $\neg \text{collinear } \{a, b, c\}$
by (*metis* *H assms*(1,2) *abc affine-hull-convex-hull collinear-affine-hull-collinear*)
moreover then have $\text{distinct } [a, b, c]$
by (*metis* *collinear-2 distinct.simps*(2) *distinct-singleton empty-set insert-absorb list.simps*(15))
moreover have $d \notin \{a, b, c\}$ **using** *d* **by** *blast*
ultimately have *?thesis*
using *abc d convex-hull-has-vertex-split-helper*[of *?p a b c d*]
by (*metis* (*no-types, lifting*) *insert-subset interior subset-trans ep*)
}
ultimately have *?thesis* **by** *fast*
} **moreover**
{ **assume** *: $\text{card } ?ep \geq 4$
moreover have $\{a, b, c\} \subseteq ?ep \wedge \text{distinct } [a, b, c]$ **using** *abc* **by** *fastforce*
ultimately have $\text{path-image } (\text{linepath } a \ b) \cap \text{interior } H \neq \{\}$
 $\vee \text{path-image } (\text{linepath } b \ c) \cap \text{interior } H \neq \{\}$
 $\vee \text{path-image } (\text{linepath } c \ a) \cap \text{interior } H \neq \{\}$
using *convex-hull-extreme-points-vertex-split*[OF *assms*(1) *assms*(4) *] **by**
presburger
then have *?thesis*
by (*metis* (*no-types, lifting*) *ep abc insert-subset subset-trans*)
}
ultimately show *?thesis* **using** *card-ep* **by** *fastforce*
qed

lemma *convex-polygon-has-good-linepath-helper*:

assumes *polygon-of* *p vts*
assumes *convex* (*path-inside* *p* \cup *path-image* *p*)
assumes $\text{card } (\text{set } vts) > 3$
obtains *a b* **where** $\{a, b\} \subseteq \text{set } vts \wedge a \neq b \wedge \neg \text{path-image } (\text{linepath } a \ b) \subseteq \text{path-image } p$
proof–
let *?H* = *convex hull* (*set vts*)
obtain *a b* **where** *ab*: $\{a, b\} \subseteq \text{set } vts \wedge a \neq b \wedge \text{path-image } (\text{linepath } a \ b) \cap \text{interior } ?H \neq \{\}$
using *convex-hull-has-vertex-split* *assms* *polygon-vts-not-collinear* **unfolding**
polygon-of-def
by *fastforce*
moreover have *interior* *?H* = *path-inside* *p*
using *assms*(1) *assms*(2) *convex-polygon-inside-is-convex-hull-interior* *polygon-convex-iff* *polygon-of-def*
by *blast*
ultimately have $\text{path-image } (\text{linepath } a \ b) \cap \text{path-inside } p \neq \{\}$ **by** *simp*

moreover have $\text{path-inside } p \cap \text{path-image } p = \{\}$ **using** *path-inside-def* **by**
auto
moreover have $\text{path-image } (\text{linepath } a \ b) \subseteq \text{path-image } p \cup \text{path-inside } p$
by (*metis ab assms(1) assms(2) convex-polygon-is-convex-hull hull-mono path-image-linepath*
polygon-of-def segment-convex-hull sup-commute)
ultimately have $\neg \text{path-image } (\text{linepath } a \ b) \subseteq \text{path-image } p$ **by fast**
thus *?thesis* **using** *ab that* **by** *meson*
qed

lemma *convex-polygon-has-good-linepath*:
assumes *convex (path-inside p \cup path-image p)*
assumes *polygon p*
assumes $p = \text{make-polygonal-path } vts$
assumes $\text{card } (\text{set } vts) > 3$
shows $\exists a \ b. \text{good-linepath } a \ b \ vts$

proof –
let $?T = \text{convex hull } (\text{set } vts)$
have $T: \text{path-image } p \cup \text{path-inside } p = ?T$
by (*metis Un-commute assms(1) assms(2) assms(3) convex-polygon-is-convex-hull*)
obtain $a \ b$ **where** $ab: a \neq b \wedge \{a, b\} \subseteq \text{set } vts \wedge \neg \text{path-image } (\text{linepath } a \ b) \subseteq$
path-image p
using *convex-polygon-has-good-linepath-helper assms unfolding polygon-of-def*
by *metis*

let $?S = \text{path-image } (\text{linepath } a \ b)$

have $p\text{-is-frontier}: \text{frontier } ?T = \text{path-image } p$
using *convex-polygon-frontier-is-path-image assms polygon-of-def polygon-convex-iff*
by *blast*

have $\text{closure } ?T = ?T$ **by** (*simp add: finite-imp-compact*)
then have $?S \subseteq \text{closure } ?T$ **using** *ab* **by** (*simp add: hull-mono segment-convex-hull*)
moreover have $\text{convex } ?T$ **using** *convex-convex-hull* **by** *auto*
moreover have $\text{convex } ?S$ **by** *simp*
moreover have $\text{rel-interior } ?S = \text{open-segment } a \ b$
by (*metis ab path-image-linepath rel-interior-closed-segment*)
moreover have $\text{rel-interior } ?T = \text{interior } ?T$
by (*metis p-is-frontier Diff-empty ab calculation(1) frontier-def rel-interior-nonempty-interior*)
ultimately have $\text{open-segment } a \ b \subseteq \text{interior } ?T$
using *subset-rel-interior-convex* **by** (*metis ab p-is-frontier frontier-def rel-frontier-def*)
then have $(\text{open-segment } a \ b) \cap \text{path-image } p = \{\}$
using *p-is-frontier frontier-def* **by** *auto*
then have $\text{closed-segment } a \ b \cap \text{path-image } p = \{a, b\}$
by (*metis (no-types, lifting) Int-Un-distrib2 Int-absorb2 Un-commute ab assms(3)*
closed-segment-eq-open subset-trans sup-bot.right-neutral vertices-on-path-image)
then have $\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } p = \{a, b\}$ **by** *simp*
thus *?thesis*
using *ab unfolding good-linepath-def*
by (*smt (verit, ccfv-threshold) IntI UnCI UnE T assms(3) hull-mono path-image-linepath*)

segment-convex-hull subset-iff)
qed

29 Pick's Theorem

definition *integral-inside*:

$integral-inside\ p = \{x. integral-vec\ x \wedge x \in path-inside\ p\}$

definition *integral-boundary*:

$integral-boundary\ p = \{x. integral-vec\ x \wedge x \in path-image\ p\}$

29.1 Pick's Theorem Triangle Case

definition *pick-triangle*:

$pick-triangle\ p\ a\ b\ c \longleftrightarrow$
 $p = make-triangle\ a\ b\ c$
 $\wedge all-integral\ [a, b, c]$
 $\wedge distinct\ [a, b, c]$
 $\wedge \neg collinear\ \{a, b, c\}$

definition *pick-holds*:

$pick-holds\ p \longleftrightarrow$
 $(let\ I = card\ \{x. integral-vec\ x \wedge x \in path-inside\ p\}\ in$
 $let\ B = card\ \{x. integral-vec\ x \wedge x \in path-image\ p\}\ in$
 $measure\ lebesgue\ (path-inside\ p) = I + B/2 - 1)$

lemma *pick-triangle-wlog-helper*:

assumes *pick-triangle* $p\ a\ b\ c$ **and**

$I = card\ (integral-inside\ p)$ **and**

$B = card\ (integral-boundary\ p)$ **and**

$integral-inside\ p = \{\}$ **and**

$integral-vec\ d \wedge d \in path-image\ (linepath\ a\ b) \wedge d \notin \{a, b, c\}$ **and** $d \notin \{a, b, c\}$ **and**

$ih: \bigwedge p' a' b' c'. (card\ (integral-inside\ p') + card\ (integral-boundary\ p') < I + B) \implies pick-triangle\ p'\ a'\ b'\ c' \implies pick-holds\ p'$

shows $measure\ lebesgue\ (path-inside\ p) = I + B/2 - 1$

proof –

have *polygon-p*: *polygon* p **using** *triangle-is-polygon* *assms* **unfolding** *pick-triangle* **by** *presburger*

then have *polygon-of*: *polygon-of* $p\ [a, b, c, a]$

unfolding *polygon-of-def* **using** *assms* **unfolding** *make-triangle-def* *pick-triangle* **by** *auto*

let $?p' = make-polygonal-path\ [a, d, b, c, a]$

have *good-linepath* $c\ d\ [a, d, b, c, a] \wedge path-image\ (make-polygonal-path\ [a, d, b, c, a]) = path-image\ p$

using *pick-triangle-basic-split* *assms* **unfolding** *pick-triangle* **by** *presburger*

then have $*$: *good-linepath* $d\ c\ [a, d, b, c, a] \wedge path-image\ (make-polygonal-path$

```

[a, d, b, c, a]) = path-image p
  using good-linepath-comm by blast
  have polygon-new: polygon (make-polygonal-path [a, d, b, c, a])
    using polygon-linepath-split-is-polygon[OF polygon-of, of 0 a b d [a, d, b, c, a]]
  assms
  by force
  have h1: make-polygonal-path [a, d, b, c, a] = make-polygonal-path ([a, d, b, c]
@ [[a, d, b, c] ! 0])
  by auto
  have h2: good-linepath d c ([a, d, b, c] @ [[a, d, b, c] ! 0])
    using * by auto
  have h3: (1::nat) < length [a, d, b, c] ∧ (3::nat) < length [a, d, b, c]
    by auto
  then have polygon-split: is-polygon-split [a, d, b, c] 1 3
    using good-linepath-implies-polygon-split[OF polygon-new h1 h2 h3] by auto
  let ?p1 = make-polygonal-path (d # [b] @ [c, d])
  let ?p2 = make-polygonal-path ([a] @ [d, c] @ [] @ [[a, d, b, c] ! 0])
  let ?I1 = card {x. integral-vec x ∧ x ∈ path-inside ?p1}
  let ?B1 = card {x. integral-vec x ∧ x ∈ path-image ?p1}
  let ?I2 = card {x. integral-vec x ∧ x ∈ path-inside ?p2}
  let ?B2 = card {x. integral-vec x ∧ x ∈ path-image ?p2}
  have p1-triangle: ?p1 = make-triangle d b c
    unfolding make-triangle-def by auto
  have p2-triangle: ?p2 = make-triangle a d c
    unfolding make-triangle-def by auto
  have I-is: I = card {x. integral-vec x ∧ x ∈ path-inside (make-polygonal-path [a,
d, b, c, a])}
    using path-image-linepath-split[of 0 [a, b, c, a] d] * assms path-inside-def
integral-inside by presburger
  have B-is: B = card {x. integral-vec x ∧ x ∈ path-image (make-polygonal-path
[a, d, b, c, a])}
    using path-image-linepath-split[of 0 [a, b, c, a] d]
    using * assms path-inside-def integral-boundary by presburger
  have all-integral-assump: all-integral [a, d, b, c]
    using assms unfolding all-integral-def pick-triangle by force

  have dist-indh1: distinct [d, b, c]
    using assms unfolding pick-triangle by auto
  have coll-indh1: ¬ collinear {d, b, c}
    using assms pick-triangle
  by (smt (verit) collinear-3-trans dist-indh1 distinct-length-2-or-more in-path-image-imp-collinear
insert-commute)
  have path-inside-inside: path-inside (make-polygonal-path (d # [b] @ [c, d])) ⊆
path-inside p
    using polygon-split unfolding is-polygon-split-def
    by (smt (z3) * One-nat-def Un-iff append-Cons append-Nil diff-Suc-1 drop0
drop-Suc-Cons nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 path-inside-def subsetI take-Suc-Cons
take-eq-Nil2)

```

then have indh1-card1 : $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } (\text{make-polygonal-path } (d \# [b] @ [c, d]))\} \leq \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$
by (*metis* (*no-types*, *lifting*) *assms*(4) *integral-inside* *Collect-empty-eq* *card.empty* *le-zero-eq* *subsetD*)
have indh1-card2 : $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } (\text{make-polygonal-path } (d \# [b] @ [c, d]))\} < \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$

proof–
have path-image-union : $\text{path-image } (\text{make-polygonal-path } (d \# [b] @ [c, d])) = \text{path-image } (\text{linepath } d \ b) \cup \text{path-image } (\text{linepath } b \ c) \cup \text{path-image } (\text{linepath } c \ d)$
using *path-image-cons-union* *p1-triangle* *make-triangle-def*
by (*metis* (*no-types*, *lifting*) *inf-sup-aci*(6) *list.discI* *make-polygonal-path.simps*(3) *nth-Cons-0*)
have path-image-db : $\text{path-image } (\text{linepath } d \ b) \subseteq \text{path-image } p$
by (*metis* *assms*(5) *list.discI* *nth-Cons-0* *path-image-cons-union* *path-image-linepath-union* *polygon-of* *polygon-of-def* *sup.cobounded2* *sup.coboundedI1*)
have path-image-bc : $\text{path-image } (\text{linepath } b \ c) \subseteq \text{path-image } p$
using *assms*(1) *linepaths-subset-make-polygonal-path-image*[of $[a, b, c, a]$ 1]
unfolding *pick-triangle* *make-triangle-def*
by *simp*
have path-image-cd1 : $\text{path-image } (\text{linepath } c \ d) - \{c, d\} \subseteq \text{path-inside } p$
using *polygon-split* **unfolding** *is-polygon-split-def*
by (*smt* (*z3*) *One-nat-def* $\langle \text{good-linepath } c \ d \ [a, d, b, c, a] \wedge \text{path-image } (\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p \rangle$ *append-Cons* *append-Nil* *insert-commute* *nth-Cons-0* *nth-Cons-Suc* *numeral-3-eq-3* *path-image-linepath* *path-inside-def* *segment-convex-hull* *sup.cobounded2*)
have path-image-cd2 : $\{c, d\} \subseteq \text{path-image } p$
using *linepaths-subset-make-polygonal-path-image* *assms*(1) **unfolding** *pick-triangle* *make-triangle-def*
by (*metis* (*no-types*, *lifting*) $\langle \text{good-linepath } c \ d \ [a, d, b, c, a] \wedge \text{path-image } (\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p \rangle$ *good-linepath-def* *subset-trans* *vertices-on-path-image*)
have $\text{path-image } (\text{linepath } c \ d) \subseteq \text{path-image } p \cup \text{path-inside } p$
using *path-image-cd1* *path-image-cd2* **by** *auto*
moreover have $\text{integral-inside } p = \{\}$ **using** *assms* **by** *force*
ultimately have path-image-cd : $\text{integral-boundary } (\text{linepath } c \ d) \subseteq \text{integral-boundary } p$ **unfolding** *integral-inside* *integral-boundary* **by** *blast*
have $a \neq d$
using *assms*(5) **by** *auto*
have $a \neq c$
using *assms*(1) **unfolding** *pick-triangle* **by** *simp*
have $a \in \text{path-image } p$
using *assms*(1) **unfolding** *pick-triangle* *make-triangle-def* **using** *vertices-on-path-image* **by** *fastforce*
have $\text{path-image } (\text{linepath } c \ d) \cap \text{path-image } p = \{c, d\}$
using * **unfolding** *good-linepath-def*
by (*smt* (*verit*, *ccfv-SIG*) *One-nat-def* *h1* *insert-commute* *is-polygon-cut-def* *is-polygon-split-def* *nth-Cons-0* *nth-Cons-Suc* *numeral-3-eq-3* *path-image-linepath* *polygon-split* *segment-convex-hull*)

```

then have a-not-in1:  $a \notin \text{path-image (linepath } c \ d)$ 
  using a-neq-c a-neq-d a-in-image by blast
have a-not-in2:  $a \notin \text{path-image (linepath } d \ b)$ 
  using Int-closed-segment assms(5) by auto
have a-not-in3:  $a \notin \text{path-image (linepath } b \ c)$ 
by (metis (no-types, lifting) assms(1) in-path-image-imp-collinear insert-commute
pick-triangle)
  then have  $a \notin \text{path-image (linepath } d \ b) \cup \text{path-image (linepath } b \ c) \cup$ 
path-image (linepath } c \ d)
  using a-not-in1 a-not-in2 a-not-in3 by simp
  then have  $a \in \text{integral-boundary } p \wedge a \notin \text{integral-boundary (make-polygonal-path$ 
[d, b, c, d])
  using path-image-union using integral-boundary a-in-image all-integral-assump
all-integral-def by auto
  then have strict-subset: integral-boundary (make-polygonal-path [d, b, c, d])  $\subset$ 
integral-boundary } p
  using path-image-union path-image-db path-image-bc path-image-cd
  unfolding integral-boundary by auto
  have integral-inside (make-polygonal-path [d, b, c, d]) = {}
  using path-inside-inside assms unfolding integral-inside by auto
  then show ?thesis using assms(2-3) strict-subset bounded-finite
  using finite-path-inside finite-path-image
  by (simp add: integral-boundary polygon-p psubset-card-mono)
qed
have fewer-points-p1: card {x. integral-vec x  $\wedge$  x  $\in$  path-inside (make-polygonal-path
(d # [b] @ [c, d]))} +
  card {x. integral-vec x  $\wedge$  x  $\in$  path-image (make-polygonal-path (d # [b] @ [c,
d]))}
   $<$  card {x. integral-vec x  $\wedge$  x  $\in$  path-inside } p  $+$ 
  card {x. integral-vec x  $\wedge$  x  $\in$  path-image } p
  using indh1-card1 indh1-card2 by linarith
have indh-1: Sigma-Algebra.measure lebesgue (path-inside ?p1) = real ?I1 + real
?B1 / 2 - 1
  using assms fewer-points-p1 p1-triangle all-integral-assump dist-indh1 coll-indh1
all-integral-def
  unfolding pick-holds pick-triangle integral-inside integral-boundary by simp

have dist-indh2: distinct [a, d, c]
  using assms unfolding pick-triangle by auto
have coll-indh2:  $\neg$  collinear {a, d, c}
  using assms pick-triangle
  by (smt (verit) collinear-3-trans dist-indh2 distinct-length-2-or-more in-path-image-imp-collinear
insert-commute)
  have path-inside-inside: path-inside (make-polygonal-path (a # [d] @ [c, a]))  $\subseteq$ 
path-inside } p
  using polygon-split unfolding is-polygon-split-def
  by (smt (z3) * One-nat-def Un-iff append-Cons append-Nil diff-Suc-1 drop0
drop-Suc-Cons nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 path-inside-def subsetI take-Suc-Cons)

```

take-eq-Nil2)
then have *indh2-card1*: $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } (\text{make-polygonal-path } (a \# [d] @ [c, a]))\} \leq \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$
by (*metis* (*no-types*, *lifting*) *assms*(4) *integral-inside Collect-empty-eq card.empty le-zero-eq subsetD*)
have *indh2-card2*: $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } (\text{make-polygonal-path } (a \# [d] @ [c, a]))\} < \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$

proof–
have *path-image-union*: $\text{path-image } (\text{make-polygonal-path } (a \# [d] @ [c, a])) = \text{path-image } (\text{linepath } a \ d) \cup \text{path-image } (\text{linepath } d \ c) \cup \text{path-image } (\text{linepath } c \ a)$
using *path-image-cons-union p2-triangle make-triangle-def*
by (*metis Un-assoc append.left-neutral append-Cons list.discI make-polygonal-path.simps*(3) *nth-Cons-0*)
have *path-image-ad*: $\text{path-image } (\text{linepath } a \ d) \subseteq \text{path-image } p$
by (*metis* $\langle \text{good-linepath } c \ d \ [a, d, b, c, a] \wedge \text{path-image } (\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p \rangle$ *inf-sup-absorb le-iff-inf list.discI nth-Cons-0 path-image-cons-union*)
have *path-image-ca*: $\text{path-image } (\text{linepath } c \ a) \subseteq \text{path-image } p$
using *assms*(1) *linepaths-subset-make-polygonal-path-image*[of $[a, b, c, a]$ 2]
unfolding *pick-triangle make-triangle-def*
by *simp*
have *path-image-cd1*: $\text{path-image } (\text{linepath } d \ c) - \{c, d\} \subseteq \text{path-inside } p$
using *polygon-split unfolding is-polygon-split-def*
by (*smt* (*z3*) *One-nat-def* $\langle \text{good-linepath } c \ d \ [a, d, b, c, a] \wedge \text{path-image } (\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p \rangle$ *append-Cons append-Nil insert-commute nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 path-image-linepath path-inside-def segment-convex-hull sup.cobounded2*)
have *path-image-cd2*: $\{c, d\} \subseteq \text{path-image } p$
using *linepaths-subset-make-polygonal-path-image assms*(1) **unfolding** *pick-triangle make-triangle-def*
by (*metis* (*no-types*, *lifting*) $\langle \text{good-linepath } c \ d \ [a, d, b, c, a] \wedge \text{path-image } (\text{make-polygonal-path } [a, d, b, c, a]) = \text{path-image } p \rangle$ *good-linepath-def subset-trans vertices-on-path-image*)
have $\text{path-image } (\text{linepath } d \ c) \subseteq \text{path-image } p \cup \text{path-inside } p$
using *path-image-cd1 path-image-cd2* **by** *auto*
moreover have $\text{integral-inside } p = \{\}$ **using** *assms* **by** *force*
ultimately have *path-image-cd*: $\text{integral-boundary } (\text{linepath } d \ c) \subseteq \text{integral-boundary } p$ **unfolding** *integral-inside integral-boundary* **by** *blast*
have *b-neq-d*: $b \neq d$
using *assms*(5) **by** *auto*
have *b-neq-c*: $b \neq c$
using *assms*(1) **unfolding** *pick-triangle* **by** *simp*
have *b-in-image*: $b \in \text{path-image } p$
using *assms*(1) **unfolding** *pick-triangle make-triangle-def* **using** *vertices-on-path-image* **by** *fastforce*
have $\text{path-image } (\text{linepath } d \ c) \cap \text{path-image } p = \{d, c\}$
using * **unfolding** *good-linepath-def*
by (*smt* (*verit*, *ccfv-SIG*) *One-nat-def h1 insert-commute is-polygon-cut-def is-polygon-split-def nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 path-image-linepath*

```

polygon-split segment-convex-hull)
  then have b-not-in1:  $b \notin \text{path-image (linepath d c)}$ 
    using b-neq-c b-neq-d b-in-image by blast
  have b-not-in2:  $b \notin \text{path-image (linepath a d)}$ 
    using Int-closed-segment assms(5) by auto
  have b-not-in3:  $b \notin \text{path-image (linepath c a)}$ 
    by (metis (no-types, lifting) assms(1) in-path-image-imp-collinear insert-commute
pick-triangle)
  then have  $b \notin \text{path-image (linepath a d)} \cup \text{path-image (linepath d c)} \cup$ 
path-image (linepath c a)
    using b-not-in1 b-not-in2 b-not-in3 by simp
  then have  $b \in \text{integral-boundary } p \wedge b \notin \text{integral-boundary (make-polygonal-path$ 
[a, d, c, a])
    using path-image-union using integral-boundary b-in-image all-integral-assump
all-integral-def by auto
  then have strict-subset:  $\text{integral-boundary (make-polygonal-path [a, d, c, a])} \subset$ 
integral-boundary p
    using path-image-union path-image-ad path-image-ca path-image-cd
unfolding integral-boundary by auto
  have integral-inside  $(\text{make-polygonal-path [a, d, c, a]}) = \{\}$ 
    using path-inside-inside assms unfolding integral-inside by auto
  then show ?thesis using assms(2-3) strict-subset bounded-finite
    using finite-path-inside finite-path-image
    by (simp add: integral-boundary polygon-p psubset-card-mono)
qed
have fewer-points-p2:  $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside (make-polygonal-path$ 
([a, d, c, a])\}\} +
 $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image (make-polygonal-path ([a, d, c, a])\}\}$ 
 $< \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} +$ 
 $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
    using indh2-card1 indh2-card2 by simp
have indh-2:  $\text{Sigma-Algebra.measure lebesgue (path-inside ?p2)} = \text{real ?I2} + \text{real$ 
?B2 / 2 - 1
    using fewer-points-p2 using assms fewer-points-p2 p2-triangle all-integral-assump
dist-indh2 coll-indh2 all-integral-def
    unfolding pick-holds pick-triangle integral-inside integral-boundary by simp

have  $\text{Sigma-Algebra.measure lebesgue (path-inside ?p1)} = \text{real ?I1} + \text{real ?B1} /$ 
2 - 1  $\implies$ 
 $\text{Sigma-Algebra.measure lebesgue (path-inside ?p2)} = \text{real ?I2} + \text{real ?B2} / 2$ 
- 1  $\implies$ 
 $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside (make-polygonal-path [a, d, b, c,$ 
a])\}\} \implies
 $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image (make-polygonal-path [a, d, b,$ 
c, a])\}\} \implies
 $\text{all-integral [a, d, b, c]} \implies$ 
 $\text{Sigma-Algebra.measure lebesgue (path-inside (make-polygonal-path [a, d, b, c,$ 
a])} =

```

```

    real I + real B / 2 - 1
  using pick-split-union[OF polygon-split, of [a] [b] [] d c ?p'] by auto
  then have Sigma-Algebra.measure lebesgue (path-inside (make-polygonal-path [a,
d, b, c, a])) =
    real I + real B / 2 - 1
  using I-is B-is all-integral-assump indh-1 indh-2 by auto
  thus measure lebesgue (path-inside p) = I + B/2 - 1
  using path-image-linepath-split[of 0 [a, b, c, a] d] by (metis path-inside-def *)
qed

```

lemma pick-triangle-helper:

```

  assumes pick-triangle p a b c and
    I = card (integral-inside p) and
    B = card (integral-boundary p) and
    integral-inside p = {} and
    integral-vec d ∧ d ∉ {a, b, c} and d ∉ {a, b, c} and
    d ∈ path-image (linepath a b)
      ∨ d ∈ path-image (linepath b c)
      ∨ d ∈ path-image (linepath c a) and
    ih: ∧p' a' b' c'. (card (integral-inside p') + card (integral-boundary p') <
I + B) ⇒ pick-triangle p' a' b' c' ⇒ pick-holds p'
  shows measure lebesgue (path-inside p) = I + B/2 - 1

```

proof –

```

  { assume d ∈ path-image (linepath a b)
    then have ?thesis using pick-triangle-wlog-helper assms by blast
  } moreover
  { assume *: d ∈ path-image (linepath b c)
    let ?p' = make-polygonal-path (rotate-polygon-vertices [a, b, c, a] 1)
    let ?I' = card (integral-inside ?p')
    let ?B' = card (integral-boundary ?p')

    have p'-p: path-image ?p' = path-image p ∧ path-inside ?p' = path-inside p
      unfolding path-inside-def
      using assms(1) make-triangle-def pick-triangle polygon-vts-arb-rotation trian-
gle-is-polygon
      by auto

    have rotate-polygon-vertices [a, b, c, a] 1 = [b, c, a, b]
      unfolding rotate-polygon-vertices-def by simp
    then have pick-triangle-p': pick-triangle ?p' b c a
      using assms unfolding pick-triangle
      by (smt (verit, best) all-integral-def distinct-length-2-or-more insert-commute
list.simps(15) make-triangle-def)
    then have measure lebesgue (path-inside ?p') = ?I' + ?B'/2 - 1
      using pick-triangle-wlog-helper[of ?p' b c a ?I' ?B' d] assms
      using integral-boundary integral-inside * insert-commute pick-triangle-p' p'-p
      by auto
    moreover have ?I' = I ∧ ?B' = B using p'-p integral-boundary integral-inside
assms(2) assms(3) by presburger
  }

```

```

ultimately have ?thesis using p'-p by auto
} moreover
{ assume *: d ∈ path-image (linepath c a)
let ?p' = make-polygonal-path (rotate-polygon-vertices [a, b, c, a] 2)
let ?I' = card (integral-inside ?p')
let ?B' = card (integral-boundary ?p')

have p'-p: path-image ?p' = path-image p ∧ path-inside ?p' = path-inside p
  unfolding path-inside-def
  using assms(1) make-triangle-def pick-triangle polygon-vts-arb-rotation triangle-is-polygon
  by auto

have rotate-polygon-vertices [a, b, c, a] 1 = [b, c, a, b]
  unfolding rotate-polygon-vertices-def by simp
also have rotate-polygon-vertices ... 1 = [c, a, b, c]
  unfolding rotate-polygon-vertices-def by simp
ultimately have rotate-polygon-vertices [a, b, c, a] 2 = [c, a, b, c]
  by (metis Suc-1 arb-rotation-as-single-rotation)
then have pick-triangle-p': pick-triangle ?p' c a b
  using assms unfolding pick-triangle
  by (smt (verit, best) all-integral-def distinct-length-2-or-more insert-commute list.simps(15) make-triangle-def)
then have measure lebesgue (path-inside ?p') = ?I' + ?B'/2 - 1
  using pick-triangle-wlog-helper[of ?p' c a b ?I' ?B' d] assms
  using integral-boundary integral-inside * insert-commute pick-triangle-p' p'-p
  by auto
moreover have ?I' = I ∧ ?B' = B using p'-p integral-boundary integral-inside
  assms(2) assms(3) by presburger
ultimately have ?thesis using p'-p by auto
}
ultimately show ?thesis using assms by blast
qed

```

```

lemma triangle-3-split-helper:
  fixes a b :: 'a::euclidean-space
  assumes a ∈ frontier S
  assumes b ∈ interior S
  assumes convex S
  assumes closed S
  shows path-image (linepath a b) ∩ frontier S = {a}
proof -
  let ?L = path-image (linepath a b)
  have a ∈ S ∧ b ∈ S using assms frontier-subset-closed interior-subset by auto
  then have ?L ⊆ S
    using assms hull-minimal segment-convex-hull by (simp add: closed-segment-subset)
  then have ?L ⊆ closure S using assms(4) by auto
  moreover have convex ?L by simp
  moreover have ?L ∩ interior S ≠ {} using assms(2) by auto

```

moreover then have $\neg ?L \subseteq \text{rel-frontier } S$
by (*metis DiffE assms(2) interior-subset-rel-interior pathfinish-in-path-image pathfinish-linepath rel-frontier-def subsetD*)
ultimately have $\text{rel-interior } ?L \subseteq \text{rel-interior } S$
using *subset-rel-interior-convex[of ?L S] assms* **by** *fastforce*
then have *open-segment a b* $\subseteq \text{interior } S$
by (*metis all-not-in-conv assms(2) empty-subsetI open-segment-eq-empty' path-image-linepath rel-interior-closed-segment rel-interior-nonempty-interior*)
moreover have $?L = \text{closed-segment } a \ b$ **by** *auto*
moreover have $\text{interior } S \cap \text{frontier } S = \{\}$ **by** (*simp add: frontier-def*)
ultimately have $?L \cap \text{frontier } S \subseteq \{a, b\}$
by (*smt (verit) Diff-iff disjoint-iff inf-commute inf-le1 open-segment-def subsetD subsetI*)
moreover have $b \notin \text{frontier } S$ **by** (*simp add: assms(2) frontier-def*)
ultimately show *?thesis* **using** *assms(1)* **by** *auto*
qed

lemma *unit-triangle-interior-point-not-collinear-e1-e2:*

assumes $p = \text{make-triangle } (\text{vector } [0, 0]) (\text{vector } [1, 0]) (\text{vector } [0, 1])$
(is $p = \text{make-triangle } ?O \ ?e1 \ ?e2$ *)*

assumes $z \in \text{path-inside } p$

shows $\neg \text{collinear } \{?O, ?e1, z\}$

proof–

have $\text{path-inside } p = \text{interior } (\text{convex hull } \{?O, ?e1, ?e2\})$

by (*metis assms(1) bounded-convex-hull bounded-empty bounded-insert convex-convex-hull convex-polygon-frontier-is-path-image2 inside-frontier-eq-interior path-inside-def triangle-convex-hull triangle-is-convex triangle-is-polygon unit-triangle-vts-not-collinear*)

then have $z \in \text{interior } (\text{convex hull } \{?O, ?e1, ?e2\})$ **using** *assms* **by** *simp*

then have $z: z\$1 > 0 \wedge z\$2 > 0$

using *assms(1) assms(2) unit-triangle-interior-char make-triangle-def* **by** *blast*

have $abc: ?O\$1 = 0 \wedge ?O\$2 = 0 \wedge ?e1\$2 = 0 \wedge ?e2\$1 = 0$ **by** *simp*

show $\neg \text{collinear } \{?O, ?e1, z\}$

proof(*rule ccontr*)

assume $\neg \neg \text{collinear } \{?O, ?e1, z\}$

then have $*$: $\text{collinear } \{?O, ?e1, z\}$ **by** *blast*

then obtain $u \ c1 \ c2$ **where** $u: ?O - ?e1 = c1 *_{\mathbb{R}} u \wedge ?e1 - z = c2 *_{\mathbb{R}} u$

unfolding *collinear-def* **by** *blast*

moreover have $c1 \neq 0$

proof–

have $(?O - ?e1)\$1 = -1$ **by** *simp*

moreover have $(?O - ?e1)\$1 = (c1 *_{\mathbb{R}} u)\1 **using** u **by** *presburger*

ultimately show *?thesis* **by** *force*

qed

moreover have $(?O - ?e1)\$2 = 0$ **by** *simp*

moreover have $(?O - ?e1)\$2 = (c1 *_{\mathbb{R}} u)\2 **by** (*simp add: calculation(1)*)

ultimately have $u\$2 = 0$ **by** *auto*

thus *False*

by (*smt (verit, ccfv-threshold) u abc scaleR-eq-0-iff vector-minus-component*)

vector-scaleR-component z)
qed
qed

lemma *triangle-interior-point-not-collinear-vertices-wlog-helper:*

assumes $p = \text{make-triangle } a \ b \ c$
assumes *polygon p*
assumes $z \in \text{path-inside } p$
shows $\neg \text{collinear } \{a, b, z\}$
proof –
let $?O = (\text{vector } [0, 0])::(\text{real}^2)$
let $?e1 = (\text{vector } [1, 0])::(\text{real}^2)$
let $?e2 = (\text{vector } [0, 1])::(\text{real}^2)$
let $?M = \text{triangle-affine } a \ b \ c$
have $a: ?M \ ?O = a$
using *triangle-affine-e1-e2* **by** *blast*
have $b: ?M \ ?e1 = b$ **using** *triangle-affine-e1-e2* **by** *simp*
have $c: ?M \ ?e2 = c$ **using** *triangle-affine-e1-e2* **by** *simp*

have *abc-not-collinear*: $\neg \text{collinear } \{a, b, c\}$
using *assms polygon-vts-not-collinear unfolding make-triangle-def polygon-of-def*
by (*metis (no-types, lifting) empty-set insertCI insert-absorb insert-commute list.simps(15)*)

have $\text{convex hull } \{a, b, c\} = \text{convex hull } \{?M \ ?O, ?M \ ?e1, ?M \ ?e2\}$
using *a b c* **by** *simp*
also have $\dots = ?M \ '(\text{convex hull } \{?O, ?e1, ?e2\})$
using *calculation triangle-affine-img* **by** *blast*
also have *interior-preserve*: $\text{interior } \dots = ?M \ '(\text{interior } (\text{convex hull } \{?O, ?e1, ?e2\}))$
using *triangle-affine-preserves-interior*[of $?M \ a \ b \ c - \text{convex hull } \{?O, ?e1, ?e2\}$]
using *abc-not-collinear*
by *presburger*
finally have $z: z \in ?M \ '(\text{interior } (\text{convex hull } \{?O, ?e1, ?e2\}))$
using *assms(1) assms(2) assms(3) make-triangle-def polygon-of-def triangle-inside-is-convex-hull-interior*
by *auto*
then obtain z' **where** $z': z' \in \text{interior } (\text{convex hull } \{?O, ?e1, ?e2\}) \wedge ?M \ z' = z$ **by** *fast*
then have $\neg \text{collinear } \{?O, ?e1, z'\}$
by (*metis convex-convex-hull convex-polygon-frontier-is-path-image2 finite.intros(1) finite-imp-bounded-convex-hull finite-insert inside-frontier-eq-interior path-inside-def triangle-convex-hull triangle-is-convex triangle-is-polygon unit-triangle-interior-point-not-collinear-e1-e2 unit-triangle-vts-not-collinear*)
then have $z'\text{-notin}$: $z' \notin \text{affine hull } \{?O, ?e1\}$ **using** *affine-hull-3-imp-collinear*
by *blast*
then have $?M \ z' \notin \text{affine hull } \{?M \ ?O, ?M \ ?e1\}$
proof –

have $inj\ ?M$ **using** *triangle-affine-inj abc-not-collinear* **by** *blast*
then have $?M\ z' \notin ?M\ '(affine\ hull\ \{?O,\ ?e1\})$ **using** z' -notin **by** (*simp add:*
inj-image-mem-iff)
moreover have $?M\ '(affine\ hull\ \{?O,\ ?e1\}) = affine\ hull\ \{?M\ ?O,\ ?M\ ?e1\}$
using *triangle-affine-preserves-affine-hull[of - a b c] abc-not-collinear* **by** *simp*
ultimately show $?thesis$ **by** *blast*
qed
then have $z \notin affine\ hull\ \{a,\ b\}$ **using** $a\ b\ z'$ **by** *argo*
thus $?thesis$
by (*metis interior-preserve z affine-hull-convex-hull affine-hull-nonempty-interior*
collinear-2 collinear-3-affine-hull collinear-affine-hull-collinear empty-iff insert-absorb2
triangle-affine-img unit-triangle-vts-not-collinear z')
qed

lemma *triangle-interior-point-not-collinear-vertices:*

assumes $p = make_triangle\ a\ b\ c$
assumes *polygon* p
assumes $z \in path_inside\ p$
shows $\neg collinear\ \{a,\ b,\ z\} \wedge \neg collinear\ \{a,\ c,\ z\} \wedge \neg collinear\ \{b,\ c,\ z\}$
proof –
let $?p1 = make_triangle\ b\ c\ a$
let $?p2 = make_triangle\ c\ a\ b$
have $p1: ?p1 = make_polygonal_path\ (rotate_polygon_vertices\ [a,\ b,\ c,\ a]\ 1)$
using *assms unfolding make-triangle-def rotate-polygon-vertices-def* **by** *fast-*
force
have $p2: ?p2 = make_polygonal_path\ (rotate_polygon_vertices\ [a,\ b,\ c,\ a]\ 2)$
using *assms unfolding make-triangle-def rotate-polygon-vertices-def* **by** (*simp*
add: numeral-Bit0)
have $path_inside\ ?p1 = path_inside\ p \wedge path_inside\ ?p2 = path_inside\ p$
using $p1\ p2$ **unfolding** *path-inside-def*
using *assms(1) assms(2) make-triangle-def polygon-vts-arb-rotation* **by** *force*
then have $z \in path_inside\ ?p1 \wedge z \in path_inside\ ?p2$ **using** *assms* **by** *force*
moreover have *polygon* $?p1 \wedge ?p2$
using *assms make-triangle-def p1 p2 rotation-is-polygon* **by** *presburger*
ultimately show $?thesis$
using *assms triangle-interior-point-not-collinear-vertices-wlog-helper*
by (*smt (verit, best) insert-commute*)
qed

lemma *triangle-3-split:*

assumes $p = make_triangle\ a\ b\ c$
assumes *polygon* p
assumes $z \in path_inside\ p$
shows *is-polygon-split-path* $[a,\ b,\ c]\ 0\ 1\ [z]$
is-polygon-split $[a,\ z,\ b,\ c]\ 1\ 3$
 $a \notin path_image\ (make_triangle\ z\ b\ c) \cup path_inside\ (make_triangle\ z\ b\ c)$
 $b \notin path_image\ (make_triangle\ a\ z\ c) \cup path_inside\ (make_triangle\ a\ z\ c)$

$c \notin \text{path-image } (\text{make-triangle } a \ b \ z) \cup \text{path-inside } (\text{make-triangle } a \ b \ z)$
proof –
let $?q = \text{make-polygonal-path } [a, z, b, c, a]$
let $?cutpath = \text{make-polygonal-path } [a, z, b]$
let $?vts = [a, b, c, a]$

let $?l1 = \text{linepath } a \ z$
let $?l2 = \text{linepath } z \ b$
let $?S = \text{path-inside } p \cup \text{path-image } p$
have $\text{convex } (\text{path-inside } p)$
using $\text{triangle-is-convex } \text{assms}(1,2) \ \text{polygon-vts-not-collinear}$ **unfolding** make-triangle-def
by $(\text{simp add: polygon-of-def triangle-inside-is-convex-hull-interior})$
then have $\text{convex: convex } (\text{path-inside } p \cup \text{path-image } p)$
using $\text{polygon-convex-iff } \text{assms}(2)$ **by** simp
then have $\text{frontier: frontier } ?S = \text{path-image } p$
using $\text{convex-polygon-frontier-is-path-image3}$ **by** $(\text{simp add: assms}(2) \ \text{sup-commute})$
have $\text{interior: interior } ?S = \text{path-inside } p$
by $(\text{metis Jordan-inside-outside-real2 closed-path-def } \langle \text{convex } (\text{path-inside } p) \rangle$
 $\text{assms}(2) \ \text{closure-Un-frontier convex-interior-closure interior-open path-inside-def}$
 $\text{polygon-def})$

have $\text{not-collinear: } \neg \text{collinear } \{a, b, z\} \wedge \neg \text{collinear } \{a, c, z\} \wedge \neg \text{collinear}$
 $\{b, c, z\}$
using $\text{triangle-interior-point-not-collinear-vertices } \text{assms}(1) \ \text{assms}(2) \ \text{assms}(3)$
by blast

have $a = \text{pathstart } ?cutpath \wedge b = \text{pathfinish } ?cutpath$ **by** simp
moreover have $a \neq b$
by $(\text{metis } \text{assms}(1) \ \text{assms}(2) \ \text{constant-linepath-is-not-loop-free make-polygonal-path.simps}(4)$
 $\text{make-triangle-def not-loop-free-first-component polygon-def simple-path-def})$
moreover have $\text{polygon } p$ **by** $(\text{simp add: assms}(2))$
moreover have $\{a, b\} \subseteq \text{set } ?vts$ **by** force
moreover have $\text{simple-path } ?cutpath$
by $(\text{simp add: insert-commute not-collinear not-collinear-loopfree-path sim-}$
 $\text{ple-path-def})$
moreover have $\text{path-image } ?cutpath \cap \text{path-image } p = \{a, b\}$
proof –
have $\{a, b\} \subseteq \text{path-image } ?cutpath \cap \text{path-image } p$
by $(\text{metis } (\text{no-types, lifting}) \ \text{Int-subset-iff Un-subset-iff } \text{assms}(1) \ \text{insert-is-Un}$
 $\text{list.simps}(15) \ \text{make-triangle-def vertices-on-path-image})$
moreover have $\text{path-image } ?cutpath \cap \text{path-image } p \subseteq \{a, b\}$
proof –
have $z \in \text{interior } ?S$ **using** assms interior **by** fast
moreover then have $a \in \text{frontier } ?S \wedge b \in \text{frontier } ?S$
using $\text{vertices-on-path-image}$
using $\langle \{a, b\} \subseteq \text{path-image } (\text{make-polygonal-path } [a, z, b]) \cap \text{path-image } p \rangle$
 frontier **by** force
moreover have $\text{closed } ?S$ **using** $\text{frontier frontier-subset-eq}$ **by** auto
ultimately have $\text{path-image } ?l1 \cap \text{path-image } p = \{a\} \wedge \text{path-image } ?l2 \cap$

```

path-image p = {b}
  using triangle-3-split-helper convex frontier
  by (metis (no-types, lifting) insert-commute path-image-linepath segment-convex-hull)
  moreover have path-image ?cutpath = path-image ?l1  $\cup$  path-image ?l2
  by (metis list.discI make-polygonal-path.simps(3) nth-Cons-0 path-image-cons-union)
  ultimately show ?thesis by blast
qed
ultimately show ?thesis by blast
qed
moreover have path-image ?cutpath  $\cap$  path-inside p  $\neq$  {}
  by (metis (no-types, opaque-lifting) Int-Un-distrib2 Un-absorb2 Un-empty assms(3)
insert-disjoint(2) list.simps(15) vertices-on-path-image)
ultimately have cutpath: is-polygon-cut-path ?vts ?cutpath
  using assms unfolding make-triangle-def is-polygon-cut-path-def by simp
thus 1: is-polygon-split-path [a, b, c] 0 1 [z]
  using polygon-cut-path-to-split-path assms(2) by (simp add: assms(1,2) make-triangle-def)

let ?l = linepath z c
let ?vts = [a, z, b, c, a]

have c-noton-cutpath: c  $\notin$  path-image ?cutpath
  by (smt (verit) UnE assms(1) assms(2) assms(3) in-path-image-imp-collinear
insert-commute make-polygonal-path.simps(3) neq-Nil-conv nth-Cons-0 path-image-cons-union
triangle-interior-point-not-collinear-vertices)

have z  $\neq$  c
proof-
  have c  $\in$  path-image p
  by (metis assms(1) insert-subset list.simps(15) make-triangle-def vertices-on-path-image)
  moreover have path-image p  $\cap$  path-inside p = {}
  by (simp add: disjoint-iff inside-def path-inside-def)
  ultimately show ?thesis using assms(3) by blast
qed
moreover have polygon-q: polygon ?q
  using 1 unfolding is-polygon-split-path-def

  by (smt (z3) One-nat-def append-Cons append-Nil diff-self-eq-0 drop0 drop-append
length-Cons length-drop length-greater-0-conv list.size(3) nth-Cons-0 nth-Cons-Suc
take-0)
  moreover have {z, c}  $\subseteq$  set ?vts by force
  moreover have l-q-int: path-image ?l  $\cap$  path-image ?q = {z, c}
  proof-
    have {z, c}  $\subseteq$  path-image ?l  $\cap$  path-image ?q
    by (metis (no-types, lifting) Int-subset-iff calculation(3) dual-order.trans
hull-subset path-image-linepath segment-convex-hull vertices-on-path-image)
    moreover
    { fix x
      assume *: x  $\in$  path-image ?l  $\cap$  path-image ?q  $\wedge$  x  $\neq$  z  $\wedge$  x  $\neq$  c
      then have x  $\in$  path-image ?q by blast
    }

```

```

then have  $x \in \text{path-image } (\text{linepath } a \ z)$ 
   $\vee x \in \text{path-image } (\text{linepath } z \ b)$ 
   $\vee x \in \text{path-image } (\text{linepath } b \ c)$ 
   $\vee x \in \text{path-image } (\text{linepath } c \ a)$ 
by (metis UnE list.discI make-polygonal-path.simps(3) nth-Cons-0 path-image-cons-union)
moreover
  { assume  $x \in \text{path-image } (\text{linepath } a \ z)$ 
    then have  $x \in \text{path-image } (\text{linepath } a \ z) \wedge x \in \text{path-image } (\text{linepath } z \ c)$ 
using * by blast
    moreover have  $z \in \text{path-image } (\text{linepath } a \ z) \wedge z \in \text{path-image } (\text{linepath } z \ c)$ 
c) by simp
    moreover have  $x \neq z$  using * by blast
    ultimately have collinear  $\{a, z, c\}$ 
      by (smt (verit, best) collinear-3-trans in-path-image-imp-collinear insert-commute)
    then have False using not-collinear by (simp add: insert-commute)
  } moreover
  { assume  $x \in \text{path-image } (\text{linepath } z \ b)$ 
    then have  $x \in \text{path-image } (\text{linepath } z \ b) \wedge x \in \text{path-image } (\text{linepath } z \ c)$ 
using * by blast
    moreover have  $z \in \text{path-image } (\text{linepath } z \ b) \wedge z \in \text{path-image } (\text{linepath } z \ c)$ 
c) by simp
    moreover have  $x \neq z$  using * by blast
    ultimately have collinear  $\{z, b, c\}$ 
      by (smt (verit, best) collinear-3-trans in-path-image-imp-collinear insert-commute)
    then have False using not-collinear by (simp add: insert-commute)
  } moreover
  { assume  $x \in \text{path-image } (\text{linepath } b \ c)$ 
    then have  $x \in \text{path-image } (\text{linepath } b \ c) \wedge x \in \text{path-image } (\text{linepath } z \ c)$ 
using * by blast
    moreover have  $c \in \text{path-image } (\text{linepath } b \ c) \wedge z \in \text{path-image } (\text{linepath } z \ c)$ 
c) by simp
    moreover have  $x \neq c$  using * by blast
    ultimately have collinear  $\{b, z, c\}$ 
      by (smt (verit, best) collinear-3-trans in-path-image-imp-collinear insert-commute)
    then have False using not-collinear by (simp add: insert-commute)
  } moreover
  { assume  $x \in \text{path-image } (\text{linepath } c \ a)$ 
    then have  $x \in \text{path-image } (\text{linepath } c \ a) \wedge x \in \text{path-image } (\text{linepath } z \ c)$ 
using * by blast
    moreover have  $c \in \text{path-image } (\text{linepath } c \ a) \wedge z \in \text{path-image } (\text{linepath } z \ c)$ 
c) by simp
    moreover have  $x \neq c$  using * by blast
    ultimately have collinear  $\{a, z, c\}$ 
      by (smt (verit, best) collinear-3-trans in-path-image-imp-collinear insert-commute)
    then have False using not-collinear by (simp add: insert-commute)
  }

```

```

    }
    ultimately have False by blast
  }
  ultimately show ?thesis by blast
qed
moreover have path-image ?l ∩ path-inside ?q ≠ {}
proof(rule ccontr)
  let ?p' = make-triangle a b z

  assume  $\neg$  path-image ?l ∩ path-inside ?q ≠ {}
  then have path-image ?l ∩ path-inside ?q = {} by blast
  then have *: rel-interior (path-image ?l) ∩ path-inside ?q = {}
    by (meson disjoint-iff rel-interior-subset subset-eq)

  have path-image ?l ⊆ path-image p ∪ path-inside p
    by (metis UnCI assms(1) assms(3) empty-subsetI hull-minimal insert-subset
list.simps(15) local.convex make-triangle-def path-image-linepath segment-convex-hull
sup-commute vertices-on-path-image)
  then have path-image ?l ⊆ convex hull {a, b, c}
    by (smt (verit, best) assms(1) convex-polygon-is-convex-hull cutpath empty-set
insertCI insert-absorb insert-commute is-polygon-cut-path-def list.simps(15) local.convex
make-triangle-def sup-commute)
  then have rel-interior (path-image ?l) ⊆ interior (convex hull {a, b, c})
    by (smt (verit, ccfv-threshold) Diff-disjoint IntE IntI Un-upper1 assms(1)
assms(2) assms(3) calculation(4) closure-Un-frontier convex-polygon-is-convex-hull
convex-segment(1) dual-order.trans empty-iff empty-set insertCI insert-absorb2 in-
sert-commute interior list.simps(15) local.convex make-triangle-def path-image-linepath
rel-frontier-def rel-interior-nonempty-interior subsetD subset-rel-interior-convex)
  then have rel-interior: rel-interior (path-image ?l) ⊆ path-inside p
    by (smt (verit, best) assms(1) convex-polygon-is-convex-hull cutpath empty-set
insertCI insert-absorb insert-commute interior is-polygon-cut-path-def list.simps(15)
local.convex make-triangle-def)

  have (let vts1 = []; vts2 = [];
vts3 = [c]; x = a; y = b;
cutpath = ?cutpath; p = make-polygonal-path ([a, b, c] @ [[a, b, c] ! 0]);
p1 = make-polygonal-path (x # vts2 @ [y] @ rev [z] @ [x]);
p2 = make-polygonal-path (vts1 @ ([x] @ [z] @ [y]) @ vts3 @ [[a, b, c] !
0]);
c1 = make-polygonal-path (x # vts2 @ [y]); c2 = make-polygonal-path
(vts1 @ ([x] @ [z] @ [y]) @ vts3)
in is-polygon-cut-path ([a, b, c] @ [[a, b, c] ! 0]) ?cutpath ∧
polygon p ∧
polygon p1 ∧
polygon p2 ∧
path-inside p1 ∩ path-inside p2 = {} ∧
path-inside p1 ∪ path-inside p2 ∪ (path-image cutpath - {x, y}) =
path-inside p ∧
(path-image p1 - path-image cutpath) ∩ (path-image p2 - path-image

```

$?cutpath) = \{\}$ \wedge
 $path\text{-}image\ p = path\text{-}image\ p1 - path\text{-}image\ ?cutpath \cup (path\text{-}image\ p2 - path\text{-}image\ ?cutpath) \cup \{x, y\}$
using 1 unfolding *is-polygon-split-path-def* **by** *fastforce*
then have (*let*
 $p = make\text{-}polygonal\text{-}path\ ([a, b, c] @ [[a, b, c] ! 0]);$
 $p1 = make\text{-}polygonal\text{-}path\ (a \# [] @ [b] @ rev [z] @ [a]);$
 $p2 = make\text{-}polygonal\text{-}path\ ([] @ ([a] @ [z] @ [b]) @ [c] @ [[a, b, c] ! 0])$
in $path\text{-}inside\ p1 \cup path\text{-}inside\ p2 \cup (path\text{-}image\ ?cutpath - \{a, b\}) = path\text{-}inside\ p$
 $\wedge (path\text{-}image\ p1 - path\text{-}image\ ?cutpath) \cap (path\text{-}image\ p2 - path\text{-}image\ ?cutpath) = \{\}$)
by *meson*
moreover have $?q = make\text{-}polygonal\text{-}path\ ([] @ ([a] @ [z] @ [b]) @ [c] @ [[a, b, c] ! 0])$
by *simp*
moreover have $?p' = make\text{-}polygonal\text{-}path\ (a \# [] @ [b] @ rev [z] @ [a])$
unfolding *make-triangle-def* **by** *simp*
moreover have $p = make\text{-}polygonal\text{-}path\ ([a, b, c] @ [[a, b, c] ! 0])$
unfolding *assms make-triangle-def* **by** *auto*
ultimately have $path\text{-}inside\text{-}p: path\text{-}inside\ ?p'$
 $\cup path\text{-}inside\ ?q$
 $\cup (path\text{-}image\ ?cutpath - \{a, b\}) = path\text{-}inside\ p$
 $\wedge (path\text{-}image\ ?p' - path\text{-}image\ ?cutpath) \cap (path\text{-}image\ ?q - path\text{-}image\ ?cutpath) = \{\}$
using 1 unfolding *make-triangle-def is-polygon-split-path-def* **by** *metis*
moreover have $a \in path\text{-}image\ ?cutpath \wedge a \notin path\text{-}inside\ ?p' \cup path\text{-}inside\ ?q$
by (*metis* (*no-types, lifting*) *UnI1* $\langle a = pathstart\ (make\text{-}polygonal\text{-}path\ [a, z, b]) \wedge b = pathfinish\ (make\text{-}polygonal\text{-}path\ [a, z, b]) \rangle assms(1)\ assms(2)\ collinear\text{-}2\ insert\text{-}absorb2\ insert\text{-}commute\ path\text{-}inside\text{-}p\ pathstart\text{-}in\text{-}path\text{-}image\ triangle\text{-}interior\text{-}point\text{-}not\text{-}collinear\text{-}vertices\text{-}wlog\text{-}helper)
moreover have $b \in path\text{-}image\ ?cutpath \wedge b \notin path\text{-}inside\ ?p' \cup path\text{-}inside\ ?q$
by (*metis* *UnI1* $\langle a = pathstart\ (make\text{-}polygonal\text{-}path\ [a, z, b]) \wedge b = pathfinish\ (make\text{-}polygonal\text{-}path\ [a, z, b]) \rangle assms(1)\ assms(2)\ collinear\text{-}2\ insert\text{-}absorb2\ path\text{-}inside\text{-}p\ pathfinish\text{-}in\text{-}path\text{-}image\ triangle\text{-}interior\text{-}point\text{-}not\text{-}collinear\text{-}vertices\text{-}wlog\text{-}helper)
ultimately have $rel\text{-}interior\ (path\text{-}image\ ?l) \subseteq$
 $(path\text{-}inside\ ?p' - path\text{-}image\ ?cutpath)$
 $\cup (path\text{-}image\ ?cutpath - \{a, b\})$
using *rel-interior ** **by** *blast*
then have $rel\text{-}interior\ (path\text{-}image\ ?l) \subseteq path\text{-}inside\ ?p' \cup path\text{-}image\ ?cutpath$
by *blast*
moreover have $path\text{-}image\ ?cutpath \subseteq path\text{-}image\ ?p'$
proof–
have $path\text{-}image\ ?cutpath = path\text{-}image\ (linepath\ a\ z) \cup path\text{-}image\ (linepath\ z\ b)$
by (*metis* *list.discI make-polygonal-path.simps(3) nth-Cons-0 path-image-cons-union*)
moreover have $path\text{-}image\ (linepath\ a\ z) = path\text{-}image\ (linepath\ z\ a)$$$

\wedge *path-image* (*linepath* *z b*) = *path-image* (*linepath* *b z*)
by (*simp add: insert-commute*)
moreover have *path-image* (*linepath* *z a*) \subseteq *path-image* *?p'*
 \wedge *path-image* (*linepath* *b z*) \subseteq *path-image* *?p'*
unfolding *make-triangle-def*
by (*metis Un-commute Un-upper2 list.discI nth-Cons-0 path-image-cons-union sup.coboundedI2*)
ultimately show *?thesis* **by** *blast*
qed
ultimately have *rel-interior* (*path-image* *?l*) \subseteq *path-inside* *?p' \cup path-image* *?p'* **by** *fast*
then have *rel-interior* (*path-image* *?l*) \subseteq *convex hull* {*a, z, b*}
unfolding *make-triangle-def*
by (*simp add: insert-commute make-triangle-def not-collinear sup-commute triangle-convex-hull*)
then have *closure* (*rel-interior* (*path-image* *?l*)) \subseteq *closure* (*convex hull* {*a, z, b*})
using *closure-mono* **by** *blast*
then have *path-image* *?l* \subseteq *convex hull* {*a, z, b*} **by** (*simp add: convex-closure-rel-interior*)
then have *c*: *c* \in *path-image* *?p' \cup path-inside* *?p'*
unfolding *make-triangle-def*
by (*metis (no-types, lifting) IntE insertCI insert-commute l-q-int make-triangle-def not-collinear subsetD triangle-convex-hull*)

moreover have *c* \notin *path-image* *?p'*
proof –
have *c* \in *path-image* *?q* – *path-image* *?cutpath* **using** *c-noton-cutpath l-q-int*
by *auto*
moreover have (*path-image* *?p' – path-image* *?cutpath*) \cap (*path-image* *?q – path-image* *?cutpath*) = {}
using *path-inside-p* **by** *fastforce*
ultimately show *?thesis* **by** *blast*
qed
moreover have *c* \notin *path-inside* *?p'*
by (*smt (verit, ccfv-threshold) DiffI IntD1 UnI1 UnI2 \langle path-image (make-polygonal-path [a, z, b]) \cap path-image p = {a, b} \rangle \langle path-image (make-polygonal-path [a, z, b]) \subseteq path-image (make-triangle a b z) \rangle assms(1) assms(2) calculation(2) collinear-2 in-mono insert-absorb2 path-inside-p triangle-interior-point-not-collinear-vertices*)
ultimately show *False* **by** *blast*
qed
ultimately have *cutpath: is-polygon-cut* *?vts z c*
using *assms* **unfolding** *make-triangle-def is-polygon-cut-def* **by** *blast*
thus *2: is-polygon-split* [*a, z, b, c*] *1 3*
using *polygon-cut-to-split*
by (*metis One-nat-def append-Cons append-Nil diff-Suc-1 length-Cons length-greater-0-conv lessI list.discI list.size(3) nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 polygon-cut-to-split zero-less-diff*)

let *?p1* = *make-triangle a z c*

let $?p2 = \text{make-triangle } z \ b \ c$
let $?p3 = \text{make-triangle } a \ b \ z$

have $(\text{path-image } ?p1 - \text{path-image } (\text{linepath } z \ c)) \cap (\text{path-image } ?p2 - \text{path-image } (\text{linepath } z \ c)) = \{\}$
using 2 **unfolding** *make-triangle-def is-polygon-split-def*
by (*smt (z3) Int-commute One-nat-def Suc-1 append-Cons append-Nil diff-numeral-Suc diff-zero drop0 drop-Suc-Cons nth-Cons-0 nth-Cons-Suc nth-Cons-numeral pred-numeral-simps(3) take0 take-Cons-numeral take-Suc-Cons*)
moreover have $a \notin \text{path-image } (\text{linepath } z \ c) \wedge b \notin \text{path-image } (\text{linepath } z \ c)$
by (*metis (no-types, lifting) assms(1) assms(2) assms(3) in-path-image-imp-collinear insert-commute triangle-interior-point-not-collinear-vertices*)
moreover have $a \in \text{path-image } ?p1 \wedge b \in \text{path-image } ?p2$
by (*metis insert-subset list.simps(15) make-triangle-def vertices-on-path-image*)
ultimately have $a \notin \text{path-image } ?p2 \wedge b \notin \text{path-image } ?p1$ **by** *auto*
moreover have $a \notin \text{path-inside } ?p2 \wedge b \notin \text{path-inside } ?p1$
proof–
have $a \notin \text{path-inside } p$
by (*metis (no-types, lifting) assms(1) assms(2) collinear-2 insertCI insert-absorb triangle-interior-point-not-collinear-vertices*)
moreover have $b \notin \text{path-inside } p$
using *assms(1) assms(2) triangle-interior-point-not-collinear-vertices-wlog-helper*
by *fastforce*
moreover have $\text{path-inside } ?p2 \subseteq \text{path-inside } ?q$
using 2 **unfolding** *is-polygon-split-def*
by (*smt (z3) One-nat-def UnCI append-Cons diff-Suc-1 drop0 drop-Suc-Cons make-triangle-def nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 self-append-conv2 subsetI take0 take-Suc-Cons*)
moreover have $\text{path-inside } ?p1 \subseteq \text{path-inside } ?q$
using 2 **unfolding** *is-polygon-split-def*
by (*smt (z3) One-nat-def Un-assoc append-Cons diff-Suc-1 drop0 drop-Suc-Cons inf-sup-absorb le-iff-inf make-triangle-def nth-Cons-0 nth-Cons-Suc numeral-3-eq-3 self-append-conv2 sup-commute take0 take-Suc-Cons*)
moreover have $\text{path-inside } ?q \subseteq \text{path-inside } p$
using 1 **unfolding** *is-polygon-split-path-def*
by (*smt (z3) One-nat-def Un-subset-iff Un-upper1 append-Cons append-Nil assms(1) diff-zero drop0 drop-Suc-Cons make-triangle-def nth-Cons-0 nth-Cons-Suc take0*)
ultimately show *?thesis* **by** *blast*
qed
moreover show $a \notin \text{path-image } ?p2 \cup \text{path-inside } ?p2$ **using** *calculation* **by** *simp*
ultimately show $b \notin \text{path-image } ?p1 \cup \text{path-inside } ?p1$ **by** *simp*

have $(\text{path-image } ?p3 - \text{path-image } ?\text{cutpath}) \cap (\text{path-image } ?q - \text{path-image } ?\text{cutpath}) = \{\}$
using 1 **unfolding** *make-triangle-def is-polygon-split-path-def*
by (*smt (z3) One-nat-def append-Cons append-Nil diff-self-eq-0 diff-zero drop0 drop-Suc-Cons nth-Cons-0 nth-Cons-Suc rev-singleton-conv take-0*)

moreover have $c \in \text{path-image } ?q$ **using** $l\text{-}q\text{-int}$ **by** auto
ultimately have $c \notin \text{path-image } ?p3$ **using** $c\text{-noton-cutpath}$ **by** blast
moreover have $c \notin \text{path-inside } ?p3$
proof –
have $c \notin \text{path-inside } p$
using $\text{assms}(1)$ $\text{assms}(2)$ $\text{triangle-interior-point-not-collinear-vertices}$ **by**
 fastforce
moreover have $\text{path-inside } ?p3 \subseteq \text{path-inside } p$
using 1 **unfolding** $\text{is-polygon-split-path-def}$
by $(\text{smt } (z3) \text{ One-nat-def Un-assoc Un-upper1 append-Cons append-Nil}$
 $\text{assms}(1) \text{ diff-Suc-Suc diff-zero make-triangle-def nth-Cons-0 nth-Cons-Suc rev-singleton-conv}$
 $\text{take0})$
ultimately show $?thesis$ **by** blast
qed
ultimately show $c \notin \text{path-image } ?p3 \cup \text{path-inside } ?p3$ **by** blast
qed

lemma smaller-triangle :

assumes $\neg \text{collinear } \{a, b, c\} \wedge \neg \text{collinear } \{a', b', c'\}$
assumes $p = \text{make-triangle } a \ b \ c$
assumes $p' = \text{make-triangle } a' \ b' \ c'$
assumes $\text{path-inside } p \subseteq \text{path-inside } p'$
assumes $\exists d. \text{integral-vec } d \wedge d \in \text{path-image } p' \cup \text{path-inside } p' \wedge d \notin \text{path-image}$
 $p \cup \text{path-inside } p$
shows $\text{card } (\text{integral-inside } p) + \text{card } (\text{integral-boundary } p) < \text{card } (\text{integral-inside}$
 $p') + \text{card } (\text{integral-boundary } p')$

proof –

have $\text{simple-path } p$ **using** assms **unfolding** make-triangle-def
using $\text{assms}(2)$ $\text{polygon-def triangle-is-polygon}$ **by** presburger
then have $\text{finite-p}: \text{finite } (\text{integral-inside } p) \wedge \text{finite } (\text{integral-boundary } p)$ **using**
 assms **unfolding** make-triangle-def
using $\text{integral-boundary integral-inside finite-integral-points-path-image finite-integral-points-path-inside}$
by metis
have $\text{simple-path } p'$ **using** assms **unfolding** make-triangle-def
using $\text{assms}(3)$ $\text{polygon-def triangle-is-polygon}$ **by** presburger
then have $\text{finite-p}': \text{finite } (\text{integral-inside } p') \wedge \text{finite } (\text{integral-boundary } p')$ **using**
 assms **unfolding** make-triangle-def
using $\text{integral-boundary integral-inside finite-integral-points-path-image finite-integral-points-path-inside}$
by metis

have $\text{polygon } p$ **using** $\text{assms}(1,2)$ $\text{triangle-is-polygon}$ **by** blast
then have $1: (\text{integral-inside } p) \cap (\text{integral-boundary } p) = \{\}$
unfolding $\text{integral-inside integral-boundary}$ **using** $\text{inside-outside-polygon un-}$
 $\text{folding inside-outside-def}$ **by** blast

have $\text{polygon } p'$ **using** $\text{assms}(1,3)$ $\text{triangle-is-polygon}$ **by** blast
then have $2: (\text{integral-inside } p') \cap (\text{integral-boundary } p') = \{\}$
unfolding $\text{integral-inside integral-boundary}$ **using** $\text{inside-outside-polygon un-}$
 $\text{folding inside-outside-def}$ **by** blast

have *path-image-subset*: $\text{path-image } p \subseteq \text{path-image } p' \cup \text{path-inside } p'$
proof–
have *p-frontier*: $\text{path-image } p = \text{frontier } (\text{convex hull } \{a, b, c\})$
by (*simp add: assms(1) assms(2) convex-polygon-frontier-is-path-image2 triangle-convex-hull triangle-is-convex triangle-is-polygon*)
have *p'-frontier*: $\text{path-image } p' = \text{frontier } (\text{convex hull } \{a', b', c'\})$
by (*simp add: assms(1) assms(3) convex-polygon-frontier-is-path-image2 triangle-convex-hull triangle-is-convex triangle-is-polygon*)

have *p-interior*: $\text{path-inside } p = \text{interior } (\text{convex hull } \{a, b, c\})$
by (*simp add: bounded-convex-hull p-frontier inside-frontier-eq-interior path-inside-def*)
have *p'-interior*: $\text{path-inside } p' = \text{interior } (\text{convex hull } \{a', b', c'\})$
by (*simp add: bounded-convex-hull p'-frontier inside-frontier-eq-interior path-inside-def*)

have $\text{interior } (\text{convex hull } \{a, b, c\}) \subseteq \text{interior } (\text{convex hull } \{a', b', c'\})$
using *assms p-interior p'-interior* **by** *argo*
moreover have $\text{compact } (\text{convex hull } \{a, b, c\}) \wedge \text{compact } (\text{convex hull } \{a', b', c'\})$
by (*simp add: compact-convex-hull*)
ultimately have $\text{frontier } (\text{convex hull } \{a, b, c\}) \subseteq \text{interior } (\text{convex hull } \{a', b', c'\}) \cup \text{frontier } (\text{convex hull } \{a', b', c'\})$
by (*smt (verit, ccfv-threshold) Jordan-inside-outside-real2 closed-path-def <polygon p'> <polygon p> assms(1) assms(2) closure-Un closure-Un-frontier closure-convex-hull finite.emptyI finite-imp-compact finite-insert p'-frontier p'-interior p-interior path-inside-def polygon-def subset-trans sup.absorb-iff1 sup-commute triangle-convex-hull*)
then show *?thesis* **using** *p'-frontier p'-interior p-frontier* **by** *blast*
qed

have $\text{card } ((\text{integral-inside } p) \cup (\text{integral-boundary } p)) = \text{card } (\text{integral-inside } p) + \text{card } (\text{integral-boundary } p)$
using *1 finite-p* **by** (*simp add: card-Un-disjoint*)
moreover have $\text{card } ((\text{integral-inside } p') \cup (\text{integral-boundary } p')) = \text{card } (\text{integral-inside } p') + \text{card } (\text{integral-boundary } p')$
using *2 finite-p'* **by** (*simp add: card-Un-disjoint*)
moreover have $(\text{integral-inside } p) \cup (\text{integral-boundary } p) \subseteq (\text{integral-inside } p') \cup (\text{integral-boundary } p')$
using *assms path-image-subset unfolding integral-inside integral-boundary* **by** *blast*
moreover then have $(\text{integral-inside } p) \cup (\text{integral-boundary } p) \subset (\text{integral-inside } p') \cup (\text{integral-boundary } p')$ **using** *assms unfolding integral-inside integral-boundary* **by** *blast*
ultimately show *?thesis* **by** (*metis finite-Un finite-p' psubset-card-mono*)
qed

lemma *pick-elem-triangle*:

fixes $p :: R\text{-to-}R^2$

assumes *p-triangle*: $p = \text{make-triangle } a \ b \ c$

```

assumes elem-triangle: elem-triangle a b c
assumes  $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$  and
       $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
shows measure lebesgue (path-inside p) =  $I + B/2 - 1$ 
proof –
  have polygon-p: polygon p
    using p-triangle triangle-is-polygon elem-triangle
    unfolding elem-triangle-def by auto
  then have path-inside p  $\cap$  path-image p =  $\{\}$ 
    using inside-outside-polygon[of p] unfolding inside-outside-def
    by auto

  let ?p = polygon (make-polygonal-path [a, b, c, a])
  have a-neq-b:  $a \neq b$ 
    using elem-triangle unfolding elem-triangle-def
    by auto
  have b-neq-c:  $b \neq c$ 
    using elem-triangle unfolding elem-triangle-def
    by auto
  have a-neq-c:  $c \neq a$ 
    using elem-triangle unfolding elem-triangle-def
    using collinear-3-eq-affine-dependent by blast

  have path-image p  $\subseteq$  convex hull  $\{a, b, c\}$ 
    using triangle-path-image-subset-convex p-triangle by auto
  then have
     $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} \subseteq \{x. \text{integral-vec } x \wedge x \in \text{convex hull } \{a, b, c\}\}$ 
    by auto
  also have  $\dots = \{a, b, c\}$ 
    using elem-triangle unfolding elem-triangle-def by auto
  finally have  $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} \subseteq \{a, b, c\}$  .
  moreover have  $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} \supseteq \{a, b, c\}$ 

  by (smt (verit) Collect-mono-iff make-triangle-def  $\langle \{x. \text{integral-vec } x \wedge x \in \text{convex hull } \{a, b, c\}\} = \{a, b, c\} \rangle$  empty-set insert-subset list.simps(15) mem-Collect-eq p-triangle subsetD vertices-on-path-image)
  ultimately have  $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} = \{a, b, c\}$  by auto
  then have card-2:  $B = 3$ 
    using a-neq-b b-neq-c a-neq-c assms(4)
    by simp

  have  $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} = \{\}$ 
proof –
  have path-inside p  $\subseteq$  convex hull  $\{a, b, c\}$ 
    by (smt (verit, best) Diff-insert-absorb make-triangle-def convex-polygon-inside-is-convex-hull-interior empty-iff empty-set insert-Diff-single insert-commute interior-subset list.simps(15) p-triangle polygon-p elem-triangle elem-triangle-def triangle-is-convex)
  then have

```

$\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} \subseteq \{x. \text{integral-vec } x \wedge x \in \text{convex hull } \{a, b, c\}\}$
by auto
also have $\dots = \{a, b, c\}$
using $\langle \{x. \text{integral-vec } x \wedge x \in \text{convex hull } \{a, b, c\}\} = \{a, b, c\} \rangle$ **by auto**
finally have $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} \subseteq \{a, b, c\}$.
moreover have
 $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} \cap \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} = \{\}$
using $\langle \text{path-inside } p \cap \text{path-image } p = \{\} \rangle$ **by auto**
ultimately show *?thesis*
using $\langle \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} = \{a, b, c\} \rangle$ **by auto**
qed
then have *card-1*: $I = 0$
using *assms*(3)
by (*metis card.empty*)

have $I + B/2 - 1 = 1/2$
using *card-1 card-2 assms*
by auto
then show *?thesis*
using *elem-triangle-area-is-half*[*OF assms*(2)] *triangle-measure-convex-hull-measure-path-inside-same*[*OF assms*(1) *assms*(2)]
by auto
qed

lemma *pick-triangle-lemma*:
fixes $p :: R\text{-to-}R^2$
assumes $p = \text{make-triangle } a \ b \ c$ **and** *all-integral* $[a, b, c]$ **and** *distinct* $[a, b, c]$
and $\neg \text{collinear } \{a, b, c\}$
 $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ **and**
 $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$
shows *measure lebesgue* ($\text{path-inside } p$) $= I + B/2 - 1$
using *assms*
proof(*induction* $\text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} + \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ *arbitrary*: $p \ a \ b \ c \ I \ B$ *rule*:*less-induct*)
case *less*
have *polygon-p*: *polygon* p **using** *triangle-is-polygon*[*OF less.prem*s(4)] *less.prem*s(1)
by *simp*
then have *polygon-of*: *polygon-of* $p \ [a, b, c, a]$
unfolding *polygon-of-def* **using** *less.prem*s(1) **unfolding** *make-triangle-def* **by** *auto*

have *convex-hull-char*: $\text{convex hull } \{a, b, c\} = \text{path-inside } p \cup \text{path-image } p$
using *triangle-convex-hull*[*OF less.prem*s(1) *less.prem*s(4)] **by auto**
then have *interior-convex-hull*: $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} \cup \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\} = \{x \in \text{convex hull } \{a, b, c\}. \text{integral-vec } x\}$
by auto
have *uts-in-path-image*: $a \in \text{path-image } p \wedge b \in \text{path-image } p \wedge c \in \text{path-image } p$

p

```

using assms(1) unfolding make-triangle-def using vertices-on-path-image
by (metis (mono-tags, lifting) insertCI less.prems(1) list.simps(15) make-triangle-def
subset-code(1))
have integral-vts:  $\text{integral-vec } a \wedge \text{integral-vec } b \wedge \text{integral-vec } c$ 
using less.prems(2)
by (simp add: all-integral-def)
then have subset:  $\{a, b, c\} \subseteq \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
using vts-in-path-image integral-vts by simp
have finite-integral-on-path-im:  $\text{finite } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
using finite-integral-points-path-image triangle-is-polygon[OF less.prems(4)]
unfolding make-triangle-def polygon-def
using less.prems(1) make-triangle-def by auto
have B-3-if:  $B > 3$  if other-point-in-set:  $\{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
 $\neq \{a, b, c\}$ 
proof –
have  $\exists d. d \notin \{a, b, c\} \wedge d \in \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
using other-point-in-set subset
by blast
then obtain  $d$  where d-prop:  $d \notin \{a, b, c\} \wedge d \in \{x. \text{integral-vec } x \wedge x \in$ 
path-image  $p\}$ 
by auto
then have subset2:  $\{a, b, c, d\} \subseteq \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
using d-prop subset by auto
have distinct  $[a, b, c, d]$ 
using d-prop
using less.prems(3) by auto
then have card-is:  $\text{card } \{a, b, c, d\} = 4$ 
by simp
show ?thesis using subset2 card-is finite-integral-on-path-im
by (metis (no-types, lifting) Suc-le-eq card-mono eval-nat-numeral(2) less.prems(6)
semiring-norm(26) semiring-norm(27))
qed
{ assume *:  $I = 0$ 
have finite  $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ 
using finite-integral-points-path-inside triangle-is-polygon[OF less.prems(4)]
unfolding make-triangle-def
by (simp add: less.prems(1) make-triangle-def polygon-def)
then have empty-inside:  $\{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\} = \{\}$ 
using * less.prems(5) by auto

{ assume **:  $B = 3$ 
have  $\{x \in \text{convex hull } \{a, b, c\}. \text{integral-vec } x\} = \{a, b, c\}$ 
using * ** less.prems(5–6) B-3-if interior-convex-hull empty-inside
by blast
then have elem-triangle  $a\ b\ c$ 
unfolding elem-triangle-def using less.prems(4) integral-vts by simp
then have measure lebesgue (path-inside  $p$ ) =  $I + B/2 - 1$ 
```

```

    using pick-elem-triangle less.premis by auto
  }
  moreover
  { assume *: B > 3
    then obtain d where d: integral-vec d ∧ d ∈ path-image p ∧ d ∉ {a, b, c}
      by (smt (verit, del-insts) subset finite-integral-on-path-im less.premis(3)
        card-3-iff collinear-3-eq-affine-dependent less.premis(4) less.premis(6) less-not-refl
        mem-Collect-eq subsetI subset-antisym)
      have path-image (make-polygonal-path [a, b, c, a]) = path-image (linepath a
        b) ∪ path-image (linepath b c) ∪ path-image (linepath c a)
      by (metis (no-types, lifting) list.discI make-polygonal-path.simps(3) nth-Cons-0
        path-image-cons-union sup-assoc)
      then have d ∈ path-image (linepath a b)
        ∨ d ∈ path-image (linepath b c)
        ∨ d ∈ path-image (linepath c a)
      using d less.premis(1) unfolding make-triangle-def polygon-of-def
      by blast
      then have measure lebesgue (path-inside p) = I + B/2 - 1
      using pick-triangle-helper less.premis less.hyps empty-inside d
      unfolding pick-holds pick-triangle integral-inside integral-boundary
      apply simp by blast
    }
    ultimately have measure lebesgue (path-inside p) = I + B/2 - 1
    using B-3-if
    by (metis (no-types, lifting) card.empty card-insert-disjoint collinear-2 fi-
      nite.emptyI finite.insertI insert-absorb less.premis(4) less.premis(6) numeral-3-eq-3)
  }
  moreover
  { assume *: I > 0
    then obtain d where d-inside: integral-vec d ∧ d ∈ path-inside p
      using less.premis(5)
      by (metis (mono-tags, lifting) Collect-empty-eq add-0 canonically-ordered-monoid-add-class.lessE
        card-0-eq card-ge-0-finite)
      have a ∈ path-image p
      using vertices-on-path-image polygon-of unfolding polygon-of-def by fastforce
      then have a-inset: a ∈ path-inside p ∪ path-image p
      by fastforce
      have convex-hull-set: convex hull set [a, b, c, a] = path-inside p ∪ path-image
        p
      using convex-hull-char
      by (simp add: insert-commute)
      then have ad-linepath-inside: path-image (linepath a d) ⊆ path-inside p ∪
        path-image p
      using d-inside convex-polygon-means-linepaths-inside[OF polygon-of con-
        vex-hull-set a-inset]
      by blast
      have b ∈ path-image p
      using vertices-on-path-image polygon-of unfolding polygon-of-def by fastforce
      then have b-inset: b ∈ path-inside p ∪ path-image p

```

```

    by fastforce
  have bd-linepath-inside: path-image (linepath b d)  $\subseteq$  path-inside p  $\cup$  path-image
p
    using d-inside convex-polygon-means-linepaths-inside[OF polygon-of con-
vex-hull-set b-inset]
    by blast
  have c  $\in$  path-image p
  using vertices-on-path-image polygon-of unfolding polygon-of-def by fastforce
  then have c-inset: c  $\in$  path-inside p  $\cup$  path-image p
  by fastforce
  then have cd-linepath-inside: path-image (linepath c d)  $\subseteq$  path-inside p  $\cup$ 
path-image p
    using d-inside convex-hull-char convex-polygon-means-linepaths-inside[OF
polygon-of convex-hull-set c-inset]
    by blast

let ?p1 = make-triangle a d c
let ?p2 = make-triangle d b c
let ?p3 = make-triangle a b d

have triangle-split:
  is-polygon-split-path [a, b, c] 0 1 [d]
  is-polygon-split [a, d, b, c] 1 3
  a  $\notin$  path-image ?p2  $\cup$  path-inside ?p2
  b  $\notin$  path-image ?p1  $\cup$  path-inside ?p1
  c  $\notin$  path-image ?p3  $\cup$  path-inside ?p3
  using triangle-3-split[of p a b c d] less.premis d-inside polygon-p apply fastforce
  using triangle-3-split[of p a b c d] less.premis d-inside polygon-p apply fastforce
  using triangle-3-split[of p a b c d] less.premis d-inside polygon-p apply fastforce
  using triangle-3-split[of p a b c d] less.premis d-inside polygon-p apply fastforce
  using triangle-3-split[of p a b c d] less.premis d-inside polygon-p by fastforce

let ?q = make-polygonal-path [a, d, b, c, a]
let ?I1 = card (integral-inside ?p1)
let ?B1 = card (integral-boundary ?p1)
let ?I2 = card (integral-inside ?p2)
let ?B2 = card (integral-boundary ?p2)
let ?I3 = card (integral-inside ?p3)
let ?B3 = card (integral-boundary ?p3)
let ?Iq = card (integral-inside ?q)
let ?Bq = card (integral-boundary ?q)
have measure lebesgue (path-inside ?p1) = ?I1 + ?B1/2 - 1
proof-
  have path-inside ?p1  $\subseteq$  path-inside ?q
    using triangle-split(2) unfolding is-polygon-split-def
    by (smt (z3) One-nat-def Un-assoc Un-upper1 append-Cons append-Nil
diff-Suc-Suc diff-zero drop0 drop-Suc-Cons make-triangle-def nth-Cons-0 nth-Cons-Suc
numeral-3-eq-3 sup-commute take0 take-Suc-Cons)
  moreover have path-inside ?q  $\subseteq$  path-inside p

```

using *triangle-split(1)* **unfolding** *is-polygon-split-path-def*
by (*smt (z3) One-nat-def Un-assoc Un-subset-iff append-Cons append-Nil*
diff-zero drop0 drop-Suc-Cons less.prem(1) make-triangle-def nth-Cons-0 nth-Cons-Suc
sup.cobounded2 take0)
ultimately have $path\text{-}inside\ ?p1 \subseteq path\text{-}inside\ p$ **by** *blast*
moreover have $\neg collinear\ \{a, d, c\}$
by (*metis d-inside insert-commute less.prem(1) polygon-p triangle-interior-point-not-collinear-vertices*)
moreover have $\neg collinear\ \{a, b, c\}$ **by** (*simp add: less.prem(4)*)
moreover have *integral-vec b*
using *integral-vts* **by** *blast*
moreover have $b \in path\text{-}image\ p$
using *vts-in-path-image* **by** *auto*
ultimately have $card\ (integral\text{-}inside\ ?p1) + card\ (integral\text{-}boundary\ ?p1)$
 $< card\ (integral\text{-}inside\ p) + card\ (integral\text{-}boundary\ p)$
using *smaller-triangle[of a d c a b c ?p1 p] triangle-split(4) less.prem(1)*
less-imp-le-nat
by *blast*
thus *?thesis*
using *less.hyps[of ?p1 a d c]* **unfolding** *integral-inside integral-boundary*
using $\neg collinear\ \{a, d, c\}$ *all-integral-def d-inside integral-vts less.prem(1)*
less.prem(3) triangle-split(3) triangle-split(5)
by *fastforce*
qed
moreover have $measure\ lebesgue\ (path\text{-}inside\ ?p2) = ?I2 + ?B2/2 - 1$
proof–
have $path\text{-}inside\ ?p2 \subseteq path\text{-}inside\ ?q$
using *triangle-split(2)* **unfolding** *is-polygon-split-def*
by (*smt (z3) One-nat-def Un-assoc Un-upper1 append-Cons append-Nil*
diff-Suc-Suc diff-zero drop0 drop-Suc-Cons make-triangle-def nth-Cons-0 nth-Cons-Suc
numeral-3-eq-3 sup-commute take0 take-Suc-Cons)
moreover have $path\text{-}inside\ ?q \subseteq path\text{-}inside\ p$
using *triangle-split(1)* **unfolding** *is-polygon-split-path-def*
by (*smt (z3) One-nat-def Un-assoc Un-subset-iff append-Cons append-Nil*
diff-zero drop0 drop-Suc-Cons less.prem(1) make-triangle-def nth-Cons-0 nth-Cons-Suc
sup.cobounded2 take0)
ultimately have $path\text{-}inside\ ?p2 \subseteq path\text{-}inside\ p$ **by** *blast*
moreover have $\neg collinear\ \{d, b, c\}$
by (*metis d-inside insert-commute less.prem(1) polygon-p triangle-interior-point-not-collinear-vertices*)
moreover have $\neg collinear\ \{a, b, c\}$ **by** (*simp add: less.prem(4)*)
moreover have *integral-vec a*
using *integral-vts* **by** *blast*
moreover have $a \in path\text{-}image\ p$
using *vts-in-path-image* **by** *auto*
ultimately have $card\ (integral\text{-}inside\ ?p2) + card\ (integral\text{-}boundary\ ?p2)$
 $< card\ (integral\text{-}inside\ p) + card\ (integral\text{-}boundary\ p)$
using *smaller-triangle[of d b c a b c ?p2 p] triangle-split(3) less.prem(1)*
less-imp-le-nat
by *blast*
thus *?thesis*

```

    using less.hyps[of ?p2 d b c] unfolding integral-inside integral-boundary
    using ⟨¬ collinear {d, b, c}⟩ all-integral-def d-inside integral-vts less.prem(1)
less.prem(3) triangle-split(3) triangle-split(5)
    by fastforce
qed
moreover have measure lebesgue (path-inside ?p3) = ?I3 + ?B3/2 - 1
proof-
  have path-inside ?p3 ⊆ path-inside p
    using triangle-split(1) unfolding is-polygon-split-path-def
    by (smt (z3) One-nat-def Un-assoc Un-upper1 append-Cons append-Nil
diff-Suc-Suc diff-zero less.prem(1) make-triangle-def nth-Cons-0 nth-Cons-Suc rev-singleton-conv
take0)
  moreover have ¬ collinear {a, b, d}
  by (metis d-inside less.prem(1) polygon-p triangle-interior-point-not-collinear-vertices)
  moreover have ¬ collinear {a, b, c} by (simp add: less.prem(4))
  moreover have integral-vec c
    using integral-vts by blast
  moreover have c ∈ path-image p
    using vts-in-path-image by auto
  ultimately have card (integral-inside ?p3) + card (integral-boundary ?p3)
< card (integral-inside p) + card (integral-boundary p)
    using smaller-triangle[of a b d a b c ?p3 p] triangle-split(5) less.prem(1)
less-imp-le-nat
    by blast
  thus ?thesis
    using less.hyps[of ?p3 a b d] unfolding integral-inside integral-boundary
    using ⟨¬ collinear {a, b, d}⟩ all-integral-def d-inside integral-vts less.prem(1)
less.prem(3) triangle-split(3) triangle-split(5)
    by fastforce
qed
moreover have measure lebesgue (path-inside ?q) = ?Iq + ?Bq/2 - 1
  using pick-split-union[OF triangle-split(2),
    of [a] [b] [] d c ?q ?p2 ?p1 ?I2 ?B2 ?I1 ?B1 ?Iq ?Bq]
  using calculation
  unfolding integral-inside integral-boundary make-triangle-def
  using all-integral-def d-inside less.prem(2) by force
ultimately have ?case
  using pick-split-path-union[OF triangle-split(1),
    of [] [] [c] a b make-polygonal-path (a # [d] @ [b]) p ?p3 ?q ?I3 ?B3 ?Iq
?Bq I B]
  unfolding integral-inside integral-boundary make-triangle-def less.prem
  using less.prem(2) by force
}
ultimately show ?case by blast
qed

```

29.2 Pocket properties

definition *index-not-in-set* :: (real^2) list \Rightarrow (real^2) set \Rightarrow nat \Rightarrow bool

where $\text{index-not-in-set } vts \ A \ i \longleftrightarrow i \in \{i. i < \text{length } vts \wedge vts ! i \notin A\}$

definition $\text{min-index-not-in-set} :: (\text{real}^2) \text{ list} \Rightarrow (\text{real}^2) \text{ set} \Rightarrow \text{nat}$
where $\text{min-index-not-in-set } vts \ A = (\text{LEAST } i. \text{index-not-in-set } vts \ A \ i)$

definition $\text{nonzero-index-in-set} :: (\text{real}^2) \text{ list} \Rightarrow (\text{real}^2) \text{ set} \Rightarrow \text{nat} \Rightarrow \text{bool}$
where
 $\text{nonzero-index-in-set } vts \ A \ i \longleftrightarrow i \in \{i. 0 < i \wedge i < \text{length } vts \wedge vts ! i \in A\}$

definition $\text{min-nonzero-index-in-set} :: (\text{real}^2) \text{ list} \Rightarrow (\text{real}^2) \text{ set} \Rightarrow \text{nat}$ **where**
 $\text{min-nonzero-index-in-set } vts \ A = (\text{LEAST } i. \text{nonzero-index-in-set } vts \ A \ i)$

definition $\text{construct-pocket-0} :: (\text{real}^2) \text{ list} \Rightarrow (\text{real}^2) \text{ set} \Rightarrow (\text{real}^2) \text{ list}$ **where**
 $\text{construct-pocket-0 } vts \ A = \text{take } ((\text{min-nonzero-index-in-set } vts \ A) + 1) \ vts$

definition $\text{is-pocket-0} :: (\text{real}^2) \text{ list} \Rightarrow (\text{real}^2) \text{ list} \Rightarrow \text{bool}$ **where**
 $\text{is-pocket-0 } vts \ vts' \longleftrightarrow$
 $\text{polygon } (\text{make-polygonal-path } vts)$
 $\wedge (\exists i. vts' = \text{take } i \ vts)$
 $\wedge 3 \leq \text{length } vts' \wedge \text{length } vts' < \text{length } vts$
 $\wedge \text{hd } vts' \in \text{frontier } (\text{convex hull } (\text{set } vts)) \wedge \text{last } vts' \in \text{frontier } (\text{convex hull } (\text{set } vts))$
 $\wedge \text{set } (\text{tl } (\text{butlast } vts')) \subseteq \text{interior } (\text{convex hull } (\text{set } vts))$

definition $\text{fill-pocket-0} :: (\text{real}^2) \text{ list} \Rightarrow \text{nat} \Rightarrow (\text{real}^2) \text{ list}$ **where**
 $\text{fill-pocket-0 } vts \ i = (\text{hd } vts) \# (\text{drop } (i-1) \ vts)$

lemma $\text{min-nonzero-index-in-set-exists}$:
assumes $\text{set } (\text{tl } vts) \cap A \neq \{\}$
shows $\exists i. \text{nonzero-index-in-set } vts \ A \ i$

proof –

obtain v **where** $v : v \in A \cap \text{set } (\text{tl } vts)$ **using** assms **by** blast

then obtain i **where** $(\text{tl } vts)!i = v \wedge i < \text{length } (\text{tl } vts)$ **by** $(\text{meson } \text{IntD2 } \text{in-set-conv-nth})$

then obtain j **where** $vts!j = v \wedge 0 < j \wedge j < \text{length } vts$ **using** nth-tl **by** fastforce
thus $?thesis$ **unfolding** $\text{nonzero-index-in-set-def}$ **using** v **by** blast

qed

lemma $\text{min-nonzero-index-in-set-defined}$:

assumes $\text{set } (\text{tl } vts) \cap A \neq \{\}$

defines $i \equiv \text{min-nonzero-index-in-set } vts \ A$

shows $\text{nonzero-index-in-set } vts \ A \ i \wedge (\forall j < i. \neg \text{nonzero-index-in-set } vts \ A \ j)$

proof –

have $\exists i. \text{nonzero-index-in-set } vts \ A \ i$ **using** assms $\text{min-nonzero-index-in-set-exists}$ **by** blast

then have $\text{nonzero-index-in-set } vts \ A \ i$

using assms **unfolding** $\text{min-nonzero-index-in-set-def}$

using *LeastI-ex* **by** *blast*
moreover have $(\forall j < i. \neg \text{nonzero-index-in-set } vts \ A \ j)$
by (*metis assms(2) wellorder-Least-lemma(2) leD min-nonzero-index-in-set-def*)
ultimately show *?thesis* **by** *blast*
qed

lemma *min-index-not-in-set-exists*:

assumes $set \ vts \supset A$

shows $\exists i. \text{index-not-in-set } vts \ A \ i$

proof –

obtain v **where** $v \in set \ vts \wedge v \notin A$ **using** *assms* **by** *blast*

then obtain i **where** $i < length \ vts \wedge vts \ ! \ i \notin A$ **by** (*metis in-set-conv-nth*)

thus *?thesis* **unfolding** *index-not-in-set-def* **by** *blast*

qed

lemma *min-index-not-in-set-defined*:

assumes $set \ vts \supset A$

defines $i \equiv \text{min-index-not-in-set } vts \ A$

shows $\text{index-not-in-set } vts \ A \ i \wedge (\forall j < i. \neg \text{index-not-in-set } vts \ A \ j)$

proof –

have $\exists i. \text{index-not-in-set } vts \ A \ i$ **using** *assms min-index-not-in-set-exists* **by** *simp*

then have $\text{index-not-in-set } vts \ A \ i$

using *assms unfolding min-index-not-in-set-def*

using *LeastI-ex* **by** *blast*

moreover have $(\forall j < i. \neg \text{index-not-in-set } vts \ A \ j)$

by (*metis assms(2) wellorder-Least-lemma(2) leD min-index-not-in-set-def*)

ultimately show *?thesis* **by** *blast*

qed

lemma *min-nonzero-index-in-set-bound*:

assumes $set \ (tl \ vts) \cap A \neq \{\}$

shows $\text{min-nonzero-index-in-set } vts \ A < length \ vts$

using *min-nonzero-index-in-set-defined assms unfolding nonzero-index-in-set-def* **by** *blast*

lemma *construct-pocket-0-subset-vts*:

assumes $set \ (tl \ vts) \cap A \neq \{\}$

shows $set \ (\text{construct-pocket-0 } vts \ A) \subseteq set \ vts$

proof –

let $?i = \text{min-nonzero-index-in-set } vts \ A$

have $\text{nonzero-index-in-set } vts \ A \ ?i$ **using** *min-nonzero-index-in-set-defined assms* **by** *presburger*

then have $?i < length \ vts$ **unfolding** *nonzero-index-in-set-def* **by** *blast*

thus *?thesis* **unfolding** *construct-pocket-0-def* **by** (*simp add: set-take-subset*)

qed

lemma *min-index-not-in-set-0*:

assumes $set \ vts \supset A$

```

assumes vts!0 ∈ A
defines i ≡ min-index-not-in-set vts A
defines r ≡ i - 1
shows vts!r ∈ A
proof -
  have *: index-not-in-set vts A i ∧ (∀ j < i. ¬ index-not-in-set vts A j)
    using min-index-not-in-set-defined[of A vts, OF assms(1)] unfolding i-def by
blast
  moreover then have r < i
    unfolding r-def i-def min-index-not-in-set-def index-not-in-set-def
  by (metis (no-types, lifting) assms(2) bot-nat-0.not-eq-extremum diff-less mem-Collect-eq
zero-less-one)
  ultimately have ¬ index-not-in-set vts A r by blast
  thus ?thesis
    unfolding index-not-in-set-def using assms * index-not-in-set-def less-imp-diff-less
by force
qed

```

```

lemma construct-pocket-0-last-in-set:
  assumes set (tl vts) ∩ A ≠ {}
  assumes vts!0 ∈ A
  defines p ≡ construct-pocket-0 vts A
  shows last p ∈ A
proof -
  let ?i = min-nonzero-index-in-set vts A
  have *: nonzero-index-in-set vts A ?i using assms(1) min-nonzero-index-in-set-defined
by blast
  then have length p = min-nonzero-index-in-set vts A + 1
    unfolding p-def construct-pocket-0-def nonzero-index-in-set-def by simp
  then have last p = p! ?i
    by (metis add-diff-cancel-right' last-conv-nth length-0-conv zero-eq-add-iff-both-eq-0
zero-neq-one)
  also have ... = vts! ?i
    unfolding p-def construct-pocket-0-def by simp
  also have ... ∈ A using * unfolding nonzero-index-in-set-def by force
  finally show ?thesis .
qed

```

```

lemma construct-pocket-0-first-last-distinct:
  assumes card A ≥ 2
  assumes A ⊆ set vts
  assumes distinct (butlast vts)
  assumes hd vts = last vts
  shows hd (construct-pocket-0 vts A) ≠ last (construct-pocket-0 vts A)
proof -
  let ?n = min-nonzero-index-in-set vts A
  have set (tl vts) ∩ A ≠ {}
    by (metis (no-types, lifting) Diff-cancel Int-commute Int-insert-right-if1 Nat.le-diff-conv2
Suc-1 add-leD1 assms(1) assms(2) card.empty card-Diff-singleton inf.orderE list.collapse

```

list.sel(2) list.set(2) not-one-le-zero plus-1-eq-Suc subset-insert
then have *n-defined: nonzero-index-in-set vts A ?n \wedge ($\forall j < ?n. \neg$ nonzero-index-in-set vts A j)*
using *min-nonzero-index-in-set-defined by presburger*
obtain *a b where ab: $a \neq b \wedge \{a, b\} \subseteq A$ by (metis assms(1) card-2-iff ex-card)*
then obtain *i j where ij: $vts!i = a \wedge vts!j = b \wedge i < \text{length vts} \wedge j < \text{length vts} \wedge i \neq j$*
by *(metis (no-types, opaque-lifting) assms(2) in-set-conv-nth insert-subset subsetD)*

have *?thesis if *: ?n < length vts - 1*
proof-
have *?n > 0 using n-defined unfolding nonzero-index-in-set-def by blast*
then have *n-bound': ?n > 0 \wedge ?n < length (butlast vts) using * by fastforce*
then have *hd vts \neq vts! ?n*
by *(metis assms(3) distinct-Ex1 hd-conv-nth ij in-set-conv-nth length-0-conv length-pos-if-in-set less-nat-zero-code nth-butlast)*
moreover then have *vts! ?n \neq last vts using assms(4) by simp*
moreover have *last (construct-pocket-0 vts A) = vts! ?n*
using *n-defined*
unfolding *construct-pocket-0-def*
by *(metis Cons-nth-drop-Suc Suc-eq-plus1 n-bound' * last-snoc less-diff-conv list.sel(1) nth-butlast take-butlast take-hd-drop)*
moreover have *hd (construct-pocket-0 vts A) = hd vts*
unfolding *construct-pocket-0-def by force*
ultimately show *?thesis by presburger*
qed
moreover have *?thesis if *: ?n = length vts - 1*
proof-
have *{i, j} \subseteq {i. i < length vts \wedge vts ! i \in A} using ij ab by simp*
moreover have *i \neq 0 \vee j \neq 0 using ij by argo*
ultimately have *nonzero-index-in-set vts A i \vee nonzero-index-in-set vts A j*
unfolding *nonzero-index-in-set-def by simp*
then have *?n = i \vee ?n = j*
by *(metis n-defined Suc-diff-1 gr-implies-not-zero ij linorder-cases not-less-eq *)*
moreover then have *last (construct-pocket-0 vts A) = vts! ?n*
by *(metis Suc-eq-plus1 construct-pocket-0-def hd-drop-conv-nth ij snoc-eq-iff-butlast take-hd-drop)*
ultimately show *?thesis*
by *(metis (no-types, lifting) ij ab Suc-eq-plus1 assms(4) bot-nat-0.not-eq-extremum hd-conv-nth insert-subset last-conv-nth less-diff-conv list.size(3) mem-Collect-eq n-defined nat-neq-iff nonzero-index-in-set-def not-less-eq that)*
qed
ultimately show *?thesis using n-defined unfolding nonzero-index-in-set-def by fastforce*
qed

lemma *construct-pocket-is-pocket:*

assumes *polygon* (*make-polygonal-path vts*)
assumes *vts!0* \in *frontier* (*convex hull* (*set vts*))
assumes *vts!1* \notin *frontier* (*convex hull* (*set vts*))
shows *is-pocket-0 vts* (*construct-pocket-0 vts* (*set vts* \cap *frontier* (*convex hull* (*set vts*))))
proof –
let *?vts'* = *construct-pocket-0 vts* (*set vts* \cap *frontier* (*convex hull* (*set vts*)))
have *ex-i*: $\exists i. ?vts' = \text{take } i \text{ vts}$ **unfolding** *construct-pocket-0-def* **by** *blast*
moreover have $3 \leq \text{length } ?vts'$
by (*smt* (*verit*) *Cons-nth-drop-Suc IntI Int-iff One-nat-def Suc-1 Suc-diff-Suc*
Suc-lessI add-diff-cancel-right' add-gr-0 append-Nil2 assms(1) assms(2) assms(3)
butlast.simps(1) butlast.simps(2) butlast-conv-take calculation cancel-comm-monoid-add-class.diff-cancel
card.empty construct-pocket-0-def construct-pocket-0-first-last-distinct construct-pocket-0-last-in-set
convex-hull-two-vts-on-frontier diff-diff-cancel diff-is-0-eq diff-is-0-eq' drop0 empty-iff
empty-set have-wraparound-vertex hd-conv-nth hd-drop-conv-nth hd-take id-take-nth-drop
last.simps last-conv-nth last-drop last-in-set last-snoc leI le-add2 le-numeral-extra(4)
le-trans length-0-conv length-greater-0-conv length-take length-tl length-upt less-2-cases
less-numeral-extra(1) less-numeral-extra(3) linorder-not-less list.distinct(1) list.sel(2)
list.sel(3) list.size(3) min.absorb4 not-gr-zero not-less-eq-eq not-numeral-le-zero nth-mem
numeral-3-eq-3 plus-1-eq-Suc polygon-at-least-3-vertices polygon-at-least-3-vertices-wraparound
polygon-def pos2 rev.simps(1) self-append-conv2 simple-polygonal-path-vts-distinct
snoc-eq-iff-butlast subset-iff take-all-iff take-eq-Nil take-hd-drop)
moreover have *vts'-length*: $\text{length } ?vts' < \text{length } vts$
by (*metis* (*no-types, lifting*) *One-nat-def Suc-1 assms(1) calculation(1) calculation(2)*
construct-pocket-0-first-last-distinct convex-hull-two-vts-on-frontier have-wraparound-vertex
hd-conv-nth inf-le1 last-snoc leI le-add2 le-trans length-take min.absorb4 not-numeral-le-zero
numeral-3-eq-3 plus-1-eq-Suc polygon-at-least-3-vertices polygon-def simple-polygonal-path-vts-distinct
take-all-iff take-eq-Nil)
moreover have *hd* *?vts' \in frontier* (*convex hull* (*set vts*))
by (*metis* *assms(2) bot-nat-0.not-eq-extremum calculation(1) calculation(2)*
hd-conv-nth hd-take list.size(3) not-numeral-le-zero take-eq-Nil)
moreover have *last* *?vts' \in frontier* (*convex hull* (*set vts*))
by (*smt* (*verit, ccfv-SIG*) *Cons-nth-drop-Suc Int-iff assms(1) assms(2) card-length*
construct-pocket-0-last-in-set drop0 drop-eq-Nil empty-iff have-wraparound-vertex
last-drop last-in-set le-add2 le-trans linorder-not-less list.sel(3) list.simps(15) not-less-eq-eq
numeral-3-eq-3 plus-1-eq-Suc polygon-at-least-3-vertices snoc-eq-iff-butlast)
moreover have *set* (*tl* (*butlast ?vts'*)) \subseteq *interior* (*convex hull* (*set vts*))
proof –
let *?A* = (*set vts* \cap *frontier* (*convex hull* (*set vts*)))
let *?r* = *min-nonzero-index-in-set vts ?A*
have *nonzero-index-in-set vts ?A ?r*
 $\wedge (\forall j < \text{min-nonzero-index-in-set vts } ?A. \neg \text{nonzero-index-in-set vts } ?A \ j)$
by (*metis* *min-nonzero-index-in-set-defined IntI Nitpick.size-list-simp(2) One-nat-def*
add-leD1 assms(1) assms(2) calculation(2) calculation(3) empty-iff empty-set have-wraparound-vertex
last-in-set last-snoc last-tl less-one not-one-le-zero nth-mem numeral-3-eq-3 plus-1-eq-Suc)
then have $\forall i. (0 < i \wedge i < ?r) \longrightarrow vts!i \notin ?A$ **unfolding** *nonzero-index-in-set-def*
by force
then have $\forall i. (0 < i \wedge i < ?r) \longrightarrow vts!i \notin \text{frontier}(\text{convex hull}(\text{set } vts))$
using *calculation(3) construct-pocket-0-def* **by** *fastforce*

then have $\forall i. (0 < i \wedge i < ?r) \longrightarrow vts!i \in \text{interior} (\text{convex hull} (\text{set } vts))$
by (*smt* (*verit*, *ccfv-threshold*) *Cons-nth-drop-Suc* *DiffI* *IntI* *One-nat-def* *add-leD1* *assms(1)* *assms(2)* *calculation(2)* *calculation(3)* *closure-subset* *drop0* *dual-order.strict-trans2* *empty-iff* *frontier-def* *have-wraparound-vertex* *hull-subset* *inf.strict-coboundedI2* *inf.strict-order-iff* *last-drop* *last-in-set* *last-snoc* *length-greater-0-conv* *list.discI* *list.sel(3)* *min-nonzero-index-in-set-bound* *nth-mem* *numeral-3-eq-3* *plus-1-eq-Suc* *subset-eq*)
moreover have *tl* (*butlast* *?vts'*) = *drop 1* (*take* *?r* *vts*)
unfolding *construct-pocket-0-def*
by (*metis* *One-nat-def* *add-implies-diff* *antisym-conv2* *butlast-take* *construct-pocket-0-def* *drop-0* *drop-Suc* *linorder-le-cases* *take-all* *vts'-length*)
moreover have $\forall v \in \text{set} (\text{drop } 1 (\text{take } ?r \text{ } vts)). \exists i. 0 < i \wedge i < ?r \wedge vts!i = v$
proof
fix *v* **assume** ***: $v \in \text{set} (\text{drop } 1 (\text{take } ?r \text{ } vts))$
then obtain *i'* **where** *i'*: $(\text{drop } 1 (\text{take } ?r \text{ } vts))!i' = v \wedge i' < ?r - 1$
by (*smt* (*z3*) *Cons-nth-drop-Suc* *One-nat-def* *ex-i* *butlast-conv-take* *calculation(2)* *drop0* *hd-conv-nth* *hd-take* *index-less-size-conv* *length-drop* *length-take* *less-imp-le-nat* *linorder-not-less* *list.collapse* *list.sel(2)* *min.absorb4* *nth-index* *take-all-iff* *take-eq-Nil* *vts'-length*)
then have $(\text{take } ?r \text{ } vts)!(i' + 1) = v$
by (*metis* *** *add commute* *drop-eq-Nil* *empty-iff* *empty-set* *nle-le* *nth-drop*)
thus $\exists i. 0 < i \wedge i < ?r \wedge vts!i = v$
by (*metis* *add-gr-0* *i'* *less-diff-conv* *nth-take* *zero-less-one*)
qed
ultimately show *?thesis* **by** *fastforce*
qed
ultimately show *?thesis* **unfolding** *is-pocket-0-def* **using** *assms(1)* **by** *argo*
qed

lemma *exists-point-above-interior*:

fixes *a* :: *real^2*

assumes $a \in \text{interior} (\text{convex hull } S)$

obtains *x* **where** $x \in S \wedge x\$2 > a\2

proof–

have *False* **if** $\forall x \in S. x\$2 \leq a\2

proof–

have $S \subseteq \{x. x \cdot (\text{vector } [0, 1]) \leq a\$2\}$

proof(*rule subsetI*)

fix *x*

assume $x \in S$

then have $x\$2 \leq a\2 **using** *that* **by** *blast*

moreover have $x \cdot (\text{vector } [0, 1]) = x\$1 * 0 + x\$2 * 1$

by (*simp* *add: cart-eq-inner-axis* *e1e2-basis(3)*)

ultimately show $x \in \{x. x \cdot (\text{vector } [0, 1]) \leq a\$2\}$ **by** *simp*

qed

then have ***: $\text{convex hull } S \subseteq \{x. x \cdot (\text{vector } [0, 1]) \leq a\$2\}$

proof–

have $S \subseteq \{v. \text{vector } [0, 1] \cdot v \leq a \$ 2\}$

by (simp add: $\langle S \subseteq \{x. x \cdot \text{vector } [0, 1] \leq a \ \$ 2\} \rangle$ inner-commute)
 then have convex hull $S \subseteq \{v. \text{vector } [0, 1] \cdot v \leq a \ \$ 2\}$
 by (simp add: convex-halfspace-le hull-minimal)
 then show ?thesis
 by (simp add: inner-commute)
 qed
 moreover have $a \cdot (\text{vector } [0, 1]) = a\2 by (simp add: cart-eq-inner-axis
 e1e2-basis(3))
 moreover have frontier $\{x. x \cdot ((\text{vector } [0, 1])::(\text{real}^2)) \leq a\$2\}$
 = $\{x. x \cdot (\text{vector } [0, 1]) = a\$2\}$
 using frontier-halfspace-le[of $(\text{vector } [0, 1])::(\text{real}^2)$ a\$2]
 by (smt (verit) Collect-cong inner-commute vector-2(2) zero-index)
 ultimately have $a \in \text{frontier } \{x. x \cdot (\text{vector } [0, 1]) \leq a\$2\}$ by blast
 thus False
 by (metis (mono-tags, lifting) Diff-iff * assms frontier-def in-frontier-in-subset
 in-mono interior-subset)
 qed
 thus ?thesis using that by fastforce
 qed

lemma exists-point-above-convex-hull-interior:

fixes $S :: (\text{real}^2)$ set
 assumes $S \neq \{\}$
 assumes compact S
 obtains x where $x \in S - (\text{interior } (\text{convex hull } S)) \wedge (\forall y \in \text{interior } (\text{convex}$
 hull $S). x\$2 > y\$2)$
 proof –
 let $?H = \text{convex hull } S$
 let $?e2 = (\text{vector } [0, 1])::(\text{real}^2)$
 let $?f = (\lambda x. x\$2)::(\text{real}^2 \Rightarrow \text{real})$
 have continuous-on $\{x. \text{True}\}$?f by (simp add: continuous-on-component)
 moreover have compact $(\text{convex hull } S)$ using assms(2) compact-convex-hull
 by blast
 moreover from calculation have compact $(?f' ?H)$
 using compact-continuous-image continuous-on-subset by blast
 ultimately obtain $x \text{ max}$ where $x: x \in ?H \wedge ?f x = \text{max} \wedge (\forall y \in ?H. y\$2 \leq$
 $\text{max})$
 by (smt (verit) Collect-mono assms(1) convex-hull-eq-empty convex-hull-explicit
 continuous-attains-sup continuous-on-subset)

 have $?H \cap \{x. ?e2 \cdot x = \text{max}\} \neq \{\}$
 by (metis (mono-tags, lifting) cart-eq-inner-axis disjoint-iff e1e2-basis(3) in-
 ner-commute mem-Collect-eq x)
 moreover have $?H \cap \{x. ?e2 \cdot x = \text{max}\} = \{\}$ if $(\forall x \in S. x\$2 < \text{max})$
 proof –
 have $S \subseteq \{x. ?e2 \cdot x < \text{max}\}$
 using that by (simp add: cart-eq-inner-axis e1e2-basis(3) inner-commute
 subset-eq)
 moreover have convex $\{x. ?e2 \cdot x < \text{max}\}$ by (simp add: convex-halfspace-lt)

ultimately show *?thesis* **using** *hull-minimal* **by** *blast*
qed
ultimately have $\exists x \in S. x \geq \max$ **by** *force*
moreover have $?H \subseteq \{x. ?e \cdot x \leq \max\}$
using *x*
by (*simp add: cart-eq-inner-axis e1e2-basis(3) inner-commute subsetI*)
moreover then have $interior\ ?H \subseteq \{x. ?e \cdot x < \max\}$
by (*metis (mono-tags) convex-empty empty-iff inner-zero-left interior-halfspace-le interior-mono real-inner-1-left separating-hyperplane-set-0 vector-2(2) zero-index*)
ultimately have $x \notin interior\ ?H \wedge (\forall y \in interior\ ?H. x > y)$
by (*smt (verit) cart-eq-inner-axis e1e2-basis(3) in-mono inner-commute mem-Collect-eq x*)
thus *?thesis* **using** *that* $\langle \exists x \in S. \max \leq x \rangle x$ **by** *fastforce*
qed

lemma *flip-function*:

defines $M \equiv (vector\ [vector\ [1, 0], vector\ [0, -1]])::(real^2^2)$
defines $f \equiv \lambda v. M * v$
defines $g \equiv (\lambda v. vector\ [v \cdot 1, -v \cdot 2])::(real^2 \Rightarrow real^2)$
shows $inj\ f\ f = g$

proof –

have $det\ M = M_{11} * M_{22} - M_{12} * M_{21}$ **using** *det-2* **by** *blast*
thus $inj\ f$ **by** (*simp add: inj-matrix-vector-mult invertible-det-nz f-def M-def*)

have $\bigwedge x. f\ x = g\ x$

proof –

fix x
have $f\ x = vector\ [M_{11} * x_1 + M_{12} * x_2, M_{21} * x_1 + M_{22} * x_2]$
 $*\ x_2]$

by (*simp add: M-def f-def mat-vec-mult-2*)

also have $\dots = vector\ [x_1, -x_2]$ **by** (*simp add: M-def*)

finally show $f\ x = g\ x$ **using** *f-def g-def* **by** *blast*

qed

thus $f = g$ **by** (*simp add: f-def g-def*)

qed

lemma *exists-point-below-convex-hull-interior*:

fixes $S :: (real^2)\ set$

assumes $S \neq \{\}$

assumes *compact S*

obtains x **where** $x \in S - (interior\ (convex\ hull\ S)) \wedge (\forall y \in interior\ (convex\ hull\ S). x < y)$

proof –

let $?M = (vector\ [vector\ [1, 0], vector\ [0, -1]])::(real^2^2)$

let $?f = \lambda v. ?M * v$

let $?g = (\lambda v. vector\ [v \cdot 1, -v \cdot 2])::(real^2 \Rightarrow real^2)$

let $?H' = ?g'(convex\ hull\ S)$

let $?S' = ?g'S$

```

have interior: ?f'(interior (convex hull S)) = interior (convex hull (?f'S))
by (smt (verit, best) flip-function convex-hull-linear-image interior-injective-linear-image
matrix-vector-mul-linear)
have hull: ?H' = convex hull ?S'
proof -
  have (*v) (vector [vector [1, 0], vector [0, - 1]]) ' (convex hull S) = convex
hull ((*v) (vector [vector [1, 0], vector [0, - 1]]) ' S::(real, 2) vec set)
  by (simp add: convex-hull-linear-image)
  then show ?thesis
  by (simp add: flip-function)
qed
moreover have compact ?S'
proof -
  have continuous-on {x. True} ?f using matrix-vector-mult-linear-continuous-on
by blast
  then have continuous-on {x. True} ?g using flip-function by simp
  thus ?thesis using assms(2) compact-continuous-image continuous-on-subset
flip-function by blast
qed
moreover have ?S' ≠ {} using assms(1) by blast
ultimately obtain x' where x': x' ∈ ?S' - (interior ?H') ∧ (∀ y ∈ interior
?H'. x'$2 > y'$2)
  using exists-point-above-convex-hull-interior[of ?S'] by auto
moreover have ?S' - (interior ?H') = ?f'(S - (interior (convex hull S)))
proof -
  have ?f'(S - (interior (convex hull S))) = ?S' - ?f'(interior (convex hull S))
  by (metis (no-types, lifting) flip-function(1) flip-function(2) image-cong im-
age-set-diff)
  thus ?thesis using flip-function(2) interior hull by auto
qed
ultimately obtain x where ?g x = x' ∧ x ∈ S - interior (convex hull S)
  using flip-function by auto
moreover have (∀ y ∈ interior (convex hull S). x $ 2 < y $ 2)
proof clarify
  fix y
  assume y ∈ interior (convex hull S)
  then have (?g x)$2 > (?g y)$2
    using x' interior hull flip-function by (metis (no-types, lifting) calculation
image-eqI)
  thus x$2 < y$2 by simp
qed
ultimately show ?thesis using that by fast
qed

lemma exists-point-above-all:
fixes p q :: R-to-R2
defines H ≡ convex hull (path-image p ∪ path-image q)
assumes path p ∧ path q

```

assumes $p\{0 < .. < 1\} \subseteq \text{interior } H$
assumes $(p\ 0)\$2 = 0 \wedge (p\ 1)\$2 = 0$
assumes $\exists x \in p\{0 < .. < 1\}. x\$2 \geq 0$
obtains x **where** $x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 > y\$2)$
proof –
let $?S = \text{path-image } p \cup \text{path-image } q$
let $?H = \text{convex hull } ?S$
obtain x **where** $x: x \in ?S - (\text{interior } ?H) \wedge (\forall y \in \text{interior } ?H. x\$2 > y\$2)$
by *(metis exists-point-above-convex-hull-interior Un-empty assms(2) compact-Un compact-path-image path-image-nonempty)*
then have $x \notin p\{0 < .. < 1\}$ **using** $H\text{-def assms(3)}$ **by** *blast*
moreover have $x \in ?S$ **using** x **by** *blast*
ultimately have $x \in \text{path-image } q \vee x \in (\text{path-image } p) - p\{0 < .. < 1\}$ **by** *blast*
moreover have $\{0..1\} - \{0 < .. < 1\} = \{0::\text{real}, 1\}$ **by** *fastforce*
ultimately have $x \in \text{path-image } q \vee x \in p\{0, 1\}$
by *(smt (verit, best) image-diff-subset path-image-def subsetD)*
moreover have $x\$2 > (p\ 0)\$2 \wedge x\$2 > (p\ 1)\2
using $H\text{-def assms(3) assms(4) assms(5) x}$ **by** *fastforce*
ultimately have $x \in \text{path-image } q \wedge x\$2 > (p\ 0)\$2 \wedge x\$2 > (p\ 1)\$2 \wedge (\forall y \in p\{0 < .. < 1\}. x\$2 > y\$2)$
using $H\text{-def assms(3) x}$ **by** *auto*
moreover have $\text{path-image } p = p\{0 < .. < 1\} \cup \{p\ 0, p\ 1\}$
proof –
have $\{0 < .. < 1\} \cup \{0::\text{real}, 1\} = \{0..1\}$ **by** *force*
thus $?thesis$ **unfolding** path-image-def **by** *blast*
qed
ultimately show $?thesis$ **by** *(simp add: that)*
qed

lemma *exists-point-below-all:*
fixes $p\ q :: R\text{-to-}R^2$
defines $H \equiv \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes $\text{path } p \wedge \text{path } q$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } H$
assumes $(p\ 0)\$2 = 0 \wedge (p\ 1)\$2 = 0$
assumes $\exists x \in \text{path-image } p \cup \text{path-image } q. x\$2 < 0$
obtains x **where** $x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 < y\$2)$
proof –
let $?thesis' = \exists x. x \in \text{path-image } q \wedge (\forall y \in \text{path-image } p. x\$2 < y\$2)$
have $?thesis'$ **if** $\exists x \in \text{path-image } p. x\$2 < 0$
proof –
have $*$: $\exists x \in p\{0 < .. < 1\}. x\$2 < 0$
proof –
have $(p\ 0)\$2 = 0 \wedge (p\ 1)\$2 = 0$ **by** *(simp add: assms(4))*
thus $?thesis$
using $\text{that unfolding path-image-def}$
using $\text{atLeastAtMost-iff less-eq-real-def}$
by *fastforce*
qed

let $?S = \text{path-image } p \cup \text{path-image } q$
let $?H = \text{convex hull } ?S$
obtain x **where** $x: x \in ?S - (\text{interior } ?H) \wedge (\forall y \in \text{interior } ?H. x\$2 < y\$2)$
by (*metis exists-point-below-convex-hull-interior Un-empty assms(2) compact-Un compact-path-image path-image-nonempty*)
then have $x \notin p\{0 < .. < 1\}$ **using** $H\text{-def assms(3)}$ **by** *blast*
moreover have $x \in ?S$ **using** x **by** *blast*
ultimately have $x \in \text{path-image } q \vee x \in (\text{path-image } p) - p\{0 < .. < 1\}$ **by**
blast
moreover have $\{0..1\} - \{0 < .. < 1\} = \{0::\text{real}, 1\}$ **by** *fastforce*
ultimately have $x \in \text{path-image } q \vee x \in p\{0, 1\}$
by (*smt (verit, best) image-diff-subset path-image-def subsetD*)
moreover have $x\$2 < (p\ 0)\$2 \wedge x\$2 < (p\ 1)\2
by (*smt (verit, ccfv-SIG) * H-def assms(3) assms(4) subset-eq x*)
ultimately have $x\$2 < (p\ 0)\$2 \wedge x\$2 < (p\ 1)\$2 \wedge (\forall y \in p\{0 < .. < 1\}. x\$2 < y\$2)$
using $H\text{-def assms(3)}$ x **by** *blast*
moreover have $\text{path-image } p = p\{0 < .. < 1\} \cup \{p\ 0, p\ 1\}$
proof -
have $\{0 < .. < 1\} \cup \{0::\text{real}, 1\} = \{0..1\}$ **by** *force*
thus $?thesis$ **unfolding** path-image-def **by** *blast*
qed
ultimately have $\forall y \in \text{path-image } p. x\$2 < y\$2$ **by** *fast*
thus $?thesis$ **using** x **by** *fast*
qed
moreover then have $?thesis'$ **if** $\neg (\exists x \in \text{path-image } p. x\$2 < 0)$ **using** $assms(5)$
by *fastforce*
ultimately show $?thesis$ **using** $that$ **by** *blast*
qed

lemma *pocket-fill-line-int-aux:*

fixes $x\ y\ z :: \text{real}^2$
defines $a \equiv y\$1$
assumes $x = 0$
assumes $a > 0 \wedge y\$2 = 0$
assumes $z\$1 < 0 \vee z\$1 > a$
assumes $z\$2 = 0$
assumes $\text{convex } A \wedge \text{compact } A$
assumes $\{x, y, z\} \subseteq A$
assumes $\{x, y\} \subseteq \text{frontier } A$
shows $z \in \text{frontier } A \wedge \text{closed-segment } x\ y \subseteq \text{frontier } A$
proof(*rule disjE[OF assms(4)]*)
assume $z\$1 > a$
moreover have $xyz: x\$1 = 0 \wedge x\$2 = 0 \wedge y\$1 = a \wedge y\$2 = 0 \wedge z\$2 = 0$
by (*simp add: a-def assms(2) assms(3) assms(5)*)
ultimately have $y: y \in \text{path-image } (\text{linepath } x\ z)$ (**is** $- \in ?L$)
using *segment-horizontal assms(3)* **by** *force*
moreover have $y\text{-neq}: y \neq x \wedge y \neq z$
by (*metis a-def assms(2) assms(3) assms(4) not-less-iff-gr-or-eq zero-index*)

ultimately have $y \in \text{rel-interior } ?L$
by (*metis UnE closed-segment-eq-open closed-segment-idem insert-Diff insert-iff path-image-linepath rel-interior-closed-segment singleton-insert-inj-eq*)
moreover have $?L \subseteq A$ **using** *assms closed-segment-subset* **by** *auto*
moreover have $z \in \text{interior } A \cup \text{frontier } A$
by (*metis Diff-iff UnI1 UnI2 assms(6) calculation(2) closure-convex-hull convex-hull-eq frontier-def in-mono pathfinish-in-path-image pathfinish-linepath*)
ultimately have $z \in \text{frontier } A$
by (*metis (no-types, lifting) Int-iff UnE y y-neq assms(6) assms(8) compact-imp-closed insert-subset singletonD triangle-3-split-helper*)
moreover have *closed-segment* $x y \subseteq \text{frontier } A$
proof(*rule ccontr*)
assume $\neg \text{closed-segment } x y \subseteq \text{frontier } A$
then obtain v **where** $v \in \text{closed-segment } x y - \text{frontier } A$ **by** *blast*
moreover then have $v \in \text{closed-segment } x y \cap \text{interior } A$
by (*metis (no-types, lifting) DiffD1 DiffD2 DiffI Int-iff assms(6) assms(7) closed-segment-subset closure-convex-hull convex-hull-eq frontier-def insert-subset subsetD*)
moreover from *calculation* **have** $v \neq x \wedge v \neq y$ **using** *assms(8)* **by** *auto*
moreover from *calculation* **have** $v\$1 < a$
by (*smt (z3) DiffD1 a-def assms(2) assms(3) exhaust-2 segment-horizontal vec-eq-iff zero-index*)
moreover from *calculation* **have** $y \in \text{open-segment } v z$
by (*smt (z3) Diff-iff xyz insert-iff open-segment-def open-segment-idem path-image-linepath segment-horizontal y y-neq*)
ultimately have $y \in \text{interior } A$
by (*metis (no-types, lifting) IntD2 assms(6) assms(7) closure-convex-hull convex-hull-eq in-interior-closure-convex-segment insertI2 singletonI subsetD*)
thus *False* **using** *assms(8) frontier-def* **by** *auto*
qed
ultimately show $z \in \text{frontier } A \wedge \text{closed-segment } x y \subseteq \text{frontier } A$ **by** *blast*
next
assume $z\$1 < 0$
moreover have $xyz: x\$1 = 0 \wedge x\$2 = 0 \wedge y\$1 = a \wedge y\$2 = 0 \wedge z\$2 = 0$
by (*simp add: a-def assms(2) assms(3) assms(5)*)
ultimately have $x: x \in \text{path-image } (\text{linepath } y z)$ (**is** $- \in ?L'$)
using *segment-horizontal assms(3)* **by** *force*
moreover have $x\text{-neq}: y \neq x \wedge x \neq z$
by (*metis a-def assms(2) assms(3) assms(4) not-less-iff-gr-or-eq zero-index*)
ultimately have $x \in \text{rel-interior } ?L'$
by (*metis UnE closed-segment-eq-open closed-segment-idem insert-Diff insert-iff path-image-linepath rel-interior-closed-segment singleton-insert-inj-eq*)
moreover have $?L' \subseteq A$
proof–
have $y \in A \wedge z \in A$ **using** *assms* **by** *blast*
thus *thesis* **by** (*simp add: assms(6) closed-segment-subset*)
qed
moreover have $z \in \text{interior } A \cup \text{frontier } A$
by (*metis Diff-iff UnI1 UnI2 assms(6) calculation(2) closure-convex-hull con-*

vex-hull-eq frontier-def in-mono pathfinish-in-path-image pathfinish-linepath
ultimately have $z \in \text{frontier } A$
by (*metis (no-types, lifting) Int-iff UnE x x-neq assms(6) assms(8) compact-imp-closed insert-subset singletonD triangle-3-split-helper*)
moreover have $\text{closed-segment } x \ y \subseteq \text{frontier } A$
proof(*rule ccontr*)
assume $\neg \text{closed-segment } x \ y \subseteq \text{frontier } A$
then obtain v **where** $v \in \text{closed-segment } x \ y - \text{frontier } A$ **by** *blast*
moreover then have $v \in \text{closed-segment } x \ y \cap \text{interior } A$
by (*metis (no-types, lifting) DiffD1 DiffD2 DiffI Int-iff assms(6) assms(7) closed-segment-subset closure-convex-hull convex-hull-eq frontier-def insert-subset subsetD*)
moreover from *calculation* **have** $v \neq x \wedge v \neq y$ **using** *assms(8)* **by** *auto*
moreover from *calculation* **have** $v\$1 > 0$
by (*smt (z3) DiffD1 a-def assms(2) assms(3) exhaust-2 segment-horizontal vec-eq-iff zero-index*)
moreover from *calculation* **have** $x \in \text{open-segment } v \ z$
by (*smt (z3) Diff-iff xyz insert-iff open-segment-def open-segment-idem path-image-linepath segment-horizontal x x-neq*)
ultimately have $x \in \text{interior } A$
by (*metis (no-types, lifting) IntD2 assms(6) assms(7) closure-convex-hull convex-hull-eq in-interior-closure-convex-segment insertI2 singletonI subsetD*)
thus *False* **using** *assms(8) frontier-def* **by** *auto*
qed
ultimately show $z \in \text{frontier } A \wedge \text{closed-segment } x \ y \subseteq \text{frontier } A$ **by** *blast*
qed

lemma *axis-dist*:

fixes $a \ b :: \text{real}^2$
shows $a\$2 = b\$2 \implies \text{dist } a \ b = \text{dist } (a\$1) \ (b\$1)$ $a\$1 = b\$1 \implies \text{dist } a \ b = \text{dist } (a\$2) \ (b\$2)$
proof–
have $\text{dist } a \ b = \text{norm } (b - a)$ **by** (*metis dist-commute dist-norm*)
also have $\dots = \text{sqrt } ((b - a) \cdot (b - a))$ **using** *norm-eq-sqrt-inner* **by** *blast*
also have $\dots = \text{sqrt } ((b - a)\$1 * (b - a)\$1 + (b - a)\$2 * (b - a)\$2)$
by (*simp add: inner-vec-def sum-2*)
finally have $*$: $\text{dist } a \ b = \text{sqrt } ((b - a)\$1 * (b - a)\$1 + (b - a)\$2 * (b - a)\$2)$.
show $a\$2 = b\$2 \implies \text{dist } a \ b = \text{dist } (a\$1) \ (b\$1)$
 $a\$1 = b\$1 \implies \text{dist } a \ b = \text{dist } (a\$2) \ (b\$2)$
apply (*simp add: * dist-real-def*)
by (*simp add: * dist-real-def*)
qed

lemma *dist-bound-1*:

fixes $a \ b \ x :: \text{real}^2$
assumes $a\$2 = x\2
assumes $b \in \text{ball } x \ \varepsilon$
assumes $\varepsilon < \text{dist } a \ x$

shows $a\$1 < x\$1 \implies b\$1 > a\1 $a\$1 > x\$1 \implies b\$1 < a\1
proof –
have 1: $\text{dist } a \ x = \text{dist } (a\$1) \ (x\$1)$ **using** *axis-dist* *assms(1)* **by** *blast*
have 2: $\text{dist } (b\$1) \ (x\$1) < \varepsilon$
by (*metis* *assms(2)* *dist-commute* *dist-vec-nth-le* *mem-ball* *order-le-less-trans*)
show $a\$1 < x\$1 \implies b\$1 > a\1 $a\$1 > x\$1 \implies b\$1 < a\1
apply (*smt* (*verit*, *ccfv-threshold*) *assms(1)* *assms(3)* 1 2 *dist-norm* *real-norm-def*)
by (*smt* (*verit*, *ccfv-threshold*) *assms(1)* *assms(3)* 1 2 *dist-norm* *real-norm-def*)
qed

lemma *dist-bound-2*:
fixes $a \ b \ x :: \text{real}^2$
assumes $a\$1 = x\1
assumes $b \in \text{ball } x \ \varepsilon$
assumes $\varepsilon < \text{dist } a \ x$
shows $a\$2 < x\$2 \implies b\$2 > a\2 $a\$2 > x\$2 \implies b\$2 < a\2
proof –
have 1: $\text{dist } a \ x = \text{dist } (a\$2) \ (x\$2)$ **using** *axis-dist* *assms(1)* **by** *blast*
have 2: $\text{dist } (b\$2) \ (x\$2) < \varepsilon$
by (*metis* *assms(2)* *dist-commute* *dist-vec-nth-le* *mem-ball* *order-le-less-trans*)
show $a\$2 < x\$2 \implies b\$2 > a\2 $a\$2 > x\$2 \implies b\$2 < a\2
apply (*smt* (*verit*, *ccfv-threshold*) *assms(1)* *assms(3)* 1 2 *dist-norm* *real-norm-def*)
by (*smt* (*verit*, *ccfv-threshold*) *assms(1)* *assms(3)* 1 2 *dist-norm* *real-norm-def*)
qed

lemma *linepath-bound-1*:
fixes $x \ y :: \text{real}^2$
shows $a < x\$1 \wedge a < y\$1 \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ a < v\1
 $x\$1 < b \wedge y\$1 < b \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ v\$1 < b$
proof –
have *: $\forall v \in \text{path-image } (\text{linepath } x \ y). \ \exists u \in \{0..1\}. \ v = (1 - u) *_R x + u *_R y$
by (*simp* *add: image-iff* *linepath-def* *path-image-def*)
have 1: $\forall u \in \{0..1\}. \ a < ((1 - u) *_R x + u *_R y)\1 **if** $a < x\$1 \wedge a < y\1
proof *clarify*
fix u **assume** $u \in \{0..1::\text{real}\}$
then **have** *: $u \geq 0 \wedge 1 - u \geq 0$ **by** *simp*
then **show** $a < ((1 - u) *_R x + u *_R y)\1
by (*smt* (*z3*) *that scaleR-collapse* *scaleR-left-mono* *vector-add-component* *vector-scaleR-component*)
qed
have 2: $\forall u \in \{0..1\}. \ ((1 - u) *_R x + u *_R y)\$1 < b$ **if** $x\$1 < b \wedge y\$1 < b$
proof *clarify*
fix u **assume** $u \in \{0..1::\text{real}\}$
then **have** *: $u \geq 0 \wedge 1 - u \geq 0$ **by** *simp*
then **show** $((1 - u) *_R x + u *_R y)\$1 < b$
by (*smt* (*z3*) *that scaleR-collapse* *scaleR-left-mono* *vector-add-component* *vector-scaleR-component*)
qed

show $a < x\$1 \wedge a < y\$1 \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ a < v\1 **using**
 * 1 **by** *fastforce*
show $x\$1 < b \wedge y\$1 < b \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ v\$1 < b$ **using**
 * 2 **by** *fastforce*
qed

lemma *linepath-bound-2:*

fixes $x \ y :: \text{real}^2$
shows $a < x\$2 \wedge a < y\$2 \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ a < v\2
 $x\$2 < b \wedge y\$2 < b \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ v\$2 < b$
proof –
have *: $\forall v \in \text{path-image } (\text{linepath } x \ y). \ \exists u \in \{0..1\}. \ v = (1 - u) *_R x + u *_R y$
by (*simp add: image-iff linepath-def path-image-def*)
have 1: $\forall u \in \{0..1\}. \ a < ((1 - u) *_R x + u *_R y)\2 **if** $a < x\$2 \wedge a < y\2
proof *clarify*
fix u **assume** $u \in \{0..1::\text{real}\}$
then **have** *: $u \geq 0 \wedge 1 - u \geq 0$ **by** *simp*
then **show** $a < ((1 - u) *_R x + u *_R y)\2
by (*smt (z3) that scaleR-collapse scaleR-left-mono vector-add-component vector-scaleR-component*)
qed
have 2: $\forall u \in \{0..1\}. \ ((1 - u) *_R x + u *_R y)\$2 < b$ **if** $x\$2 < b \wedge y\$2 < b$
proof *clarify*
fix u **assume** $u \in \{0..1::\text{real}\}$
then **have** *: $u \geq 0 \wedge 1 - u \geq 0$ **by** *simp*
then **show** $((1 - u) *_R x + u *_R y)\$2 < b$
by (*smt (z3) that scaleR-collapse scaleR-left-mono vector-add-component vector-scaleR-component*)
qed
show $a < x\$2 \wedge a < y\$2 \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ a < v\2 **using**
 * 1 **by** *fastforce*
show $x\$2 < b \wedge y\$2 < b \implies \forall v \in \text{path-image } (\text{linepath } x \ y). \ v\$2 < b$ **using**
 * 2 **by** *fastforce*
qed

lemma *linepath-int-corner:*

fixes $x \ y \ z :: \text{real}^2$
assumes $x\$2 \neq y\2
assumes $y\$2 = z\2
shows $\text{path-image } (\text{linepath } x \ y) \cap \text{path-image } (\text{linepath } y \ z) = \{y\}$
 (is $\text{path-image } ?l1 \cap \text{path-image } ?l2 = \{y\}$)
proof –
have 1: $y \in \text{path-image } ?l1 \cap \text{path-image } ?l2$ **by** *simp*

have $\forall t \in \{0..1\}. \ (?l1 \ t)\$2 = y\$2 \implies t = 1$
proof *clarify*
fix $t :: \text{real}$
assume 1: $t \in \{0..1\}$

assume 2: $(?l1\ t)\$2 = y\2
have $(?l1\ t)\$2 = ((1 - t) * (x\$2) + t * (y\$2))$ **by** (*simp add: linepath-def*)
thus $t = 1$
by (*smt (verit, best) assms 2 distrib-right inner-real-def mult.commute real-inner-1-right vector-space-over-itself.scale-cancel-left*)
qed
then have $\forall t \in \{0..1\}. (?l1\ t)\$2 = y\$2 \longleftrightarrow t = 1$ **by** (*metis linepath-1'*)
moreover have $\forall t \in \{0..1\}. (?l2\ t)\$2 = y\$2$
unfolding *linepath-def*
by (*metis (no-types, lifting) assms(2) segment-degen-1 vector-add-component vector-scaleR-component*)
ultimately have 2: $path\text{-}image\ ?l1 \cap path\text{-}image\ ?l2 \subseteq \{y\}$
by (*smt (verit, best) 1 IntD1 IntD2 imageE path-defs(4) singleton-iff subsetI*)

show *?thesis* **using** 1 2 **by** *fastforce*
qed

lemma *linepath-int-vertical*:

fixes $w\ x\ y\ z :: real^2$
assumes $w\$1 \neq y\1
assumes $w\$1 = x\1
assumes $y\$1 = z\1
shows $path\text{-}image\ (linepath\ w\ x) \cap path\text{-}image\ (linepath\ y\ z) = \{\}$
using *assms segment-vertical* **by** *fastforce*

lemma *linepath-int-horizontal*:

fixes $w\ x\ y\ z :: real^2$
assumes $w\$2 \neq y\2
assumes $w\$2 = x\2
assumes $y\$2 = z\2
shows $path\text{-}image\ (linepath\ w\ x) \cap path\text{-}image\ (linepath\ y\ z) = \{\}$
using *assms segment-horizontal* **by** *fastforce*

lemma *linepath-int-columns*:

fixes $w\ x\ y\ z :: real^2$
assumes $w\$1 < y\$1 \wedge w\$1 < z\1
assumes $x\$1 < y\$1 \wedge x\$1 < z\1
shows $path\text{-}image\ (linepath\ w\ x) \cap path\text{-}image\ (linepath\ y\ z) = \{\}$
(is path-image ?l1 \cap path-image ?l2 = $\{\}$)

proof –

have $\forall t1 \in \{0..1\}. \forall t2 \in \{0..1\}. (?l2\ t2)\$1 > (?l1\ t1)\$1$
by (*smt (verit, ccfv-SIG) assms linepath-bound-1 linepath-in-path path-image-linepath*)
thus *?thesis* **by** (*smt (verit, best) disjoint-iff imageE path-image-def*)

qed

lemma *linepath-int-rows*:

fixes $w\ x\ y\ z :: real^2$
assumes $w\$2 < y\$2 \wedge w\$2 < z\2

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assumes  $x\$2 < y\$2 \wedge x\$2 < z\$2$ 
shows  $\text{path-image } (\text{linepath } w \ x) \cap \text{path-image } (\text{linepath } y \ z) = \{\}$ 
  (is  $\text{path-image } ?l1 \cap \text{path-image } ?l2 = \{\}$ )
proof -
  have  $\forall t1 \in \{0..1\}. \forall t2 \in \{0..1\}. (?l2 \ t2)\$2 > (?l1 \ t1)\$2$ 
  by (smt (verit, ccfv-SIG) assms linepath-bound-2 linepath-in-path path-image-linepath)
  thus ?thesis by (smt (verit, best) disjoint-iff imageE path-image-def)
qed

lemma horizontal-segment-at-0:
  assumes  $a > 0$ 
  shows  $\text{closed-segment } ((\text{vector } [0, 0])::(\text{real}^2)) (\text{vector } [a, 0]) = \{x. x\$2 = 0$ 
 $\wedge x\$1 \in \{0..a\}\}$ 
  (is  $?l = ?s$ )
proof -
  have  $?l \subseteq ?s$ 
  proof(rule subsetI)
    fix  $x$ 
    assume  $*$ :  $x \in ?l$ 
    then have  $x\$2 = 0$  using segment-horizontal by auto
    moreover have  $0 \leq x\$1 \wedge x\$1 \leq a$  using  $*$  assms segment-horizontal by
force
    ultimately show  $x \in ?s$  by force
  qed
  moreover have  $?s \subseteq ?l$ 
  proof(rule subsetI)
    fix  $x$ 
    assume  $*$ :  $x \in ?s$ 
    then have  $x = (x\$1 / a) *_R (\text{vector } [a, 0]) + (1 - (x\$1 / a)) *_R (\text{vector } [0,$ 
 $0])$ 
    proof-
      have  $(x\$1 / a) *_R ((\text{vector } [a, 0])::(\text{real}^2)) = \text{vector } [x\$1, 0]$ 
      using vec-scaleR-2 assms by fastforce
      moreover have  $(1 - (x\$1 / a)) *_R ((\text{vector } [0, 0])::(\text{real}^2)) = \text{vector } [0,$ 
 $0]$ 
      using vec-scaleR-2 by simp
      moreover have  $x = \text{vector } [x\$1, 0]$ 
      by (smt (verit) * exhaust-2 mem-Collect-eq vec-eq-iff vector-2(1) vector-2(2))
      ultimately show ?thesis
      by (metis add-cancel-right-right scaleR-collapse vec-scaleR-2 vector-2(2))
    qed
    moreover have  $x\$1 / a \in \{0..1\}$  using  $*$  assms by fastforce
    ultimately show  $x \in ?l$ 
    by (smt (verit, del-insts) add commute atLeastAtMost-iff mem-Collect-eq
closed-segment-def)
  qed
  ultimately show ?thesis by blast
qed

```

lemma *horizontal-segment-at-0'*:

fixes $x\ y :: \text{real}^2$

assumes $a > 0$

assumes $x\$1 = 0 \wedge x\$2 = 0 \wedge y\$1 = a \wedge y\$2 = 0$

shows $\text{closed-segment } x\ y = \{x. x\$2 = 0 \wedge x\$1 \in \{0..a\}\}$

proof –

have $x = \text{vector } [0, 0] \wedge y = \text{vector } [a, 0]$

by (*smt* (*verit*, *best*) *assms*(2) *exhaust-2* *vec-eq-iff* *vector-2*(1) *vector-2*(2))

thus *?thesis* **using** *horizontal-segment-at-0* *assms* **by** *presburger*

qed

lemma *pocket-fill-line-int-aux1*:

fixes $p\ q :: R\text{-to-}R^2$

defines $p0 \equiv \text{pathstart } p$

defines $p1 \equiv \text{pathfinish } p$

defines $q0 \equiv \text{pathstart } q$

defines $q1 \equiv \text{pathfinish } q$

defines $a \equiv p1\$1$

defines $l \equiv \text{closed-segment } p0\ p1$

assumes *simple-path* p

assumes *simple-path* q

assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$

assumes $a > 0$

assumes $\text{path-image } q \cap \{x. x\$2 = 0\} \subseteq l$

assumes $\text{path-image } p \cap \{x. x\$2 = 0\} \subseteq l$

assumes $\forall v \in \text{path-image } p. q0\$2 \leq v\$2$

assumes $\forall v \in \text{path-image } p. q1\$2 > v\$2$

shows $\text{path-image } p \cap \text{path-image } q \neq \{\}$

proof –

have $p0: p0 = 0$

by (*metis* (*mono-tags*, *opaque-lifting*) *assms*(9) *exhaust-2* *vec-eq-iff* *zero-index*)

moreover **have** $p1: p1 = \text{vector } [a, 0]$

by (*smt* (*verit*) *a-def* *assms*(9) *exhaust-2* *vec-eq-iff* *vector-2*(1) *vector-2*(2))

obtain $a\text{-}x$ **where** $a\text{-}x: \forall v \in \text{path-image } p \cup \text{path-image } q. a\text{-}x < v\1

proof –

let $?a\text{-}x = \text{Inf } ((\lambda v. v\$1) (\text{path-image } p \cup \text{path-image } q))$

have *compact* ($\text{path-image } p \cup \text{path-image } q$)

by (*simp* *add: assms*(7) *assms*(8) *compact-Un* *compact-simple-path-image*)

moreover **have** *continuous-on UNIV* $((\lambda v. v\$1)::(\text{real}^2 \Rightarrow \text{real}))$

by (*simp* *add: continuous-on-component*)

ultimately **have** $*$: *compact* $((\lambda v. v\$1) (\text{path-image } p \cup \text{path-image } q))$

by (*meson* *compact-continuous-image* *continuous-on-subset* *top-greatest*)

then **have** $\forall x \in ((\lambda v. v\$1) (\text{path-image } p \cup \text{path-image } q)). ?a\text{-}x \leq x$

by (*simp* *add: assms*(7) *assms*(8) *bounded-component-cart* *bounded-has-Inf*(1) *bounded-simple-path-image*)

thus *?thesis* **using** *that*[*of* $?a\text{-}x - 1$] **by** (*smt* (*verit*, *ccfv-SIG*) *assms*(10) *imageI*)

qed

obtain $b-x$ **where** $b-x: \forall v \in \text{path-image } p \cup \text{path-image } q. b-x > v\1
proof–
let $?b-x = \text{Sup } ((\lambda v. v\$1) \text{'(path-image } p \cup \text{path-image } q))$
have $\text{compact } (\text{path-image } p \cup \text{path-image } q)$
by (*simp add: assms(7) assms(8) compact-Un compact-simple-path-image*)
moreover have $\text{continuous-on UNIV } ((\lambda v. v\$1)::(\text{real}^2 \Rightarrow \text{real}))$
by (*simp add: continuous-on-component*)
ultimately have $*$: $\text{compact } ((\lambda v. v\$1) \text{'(path-image } p \cup \text{path-image } q))$
by (*meson compact-continuous-image continuous-on-subset top-greatest*)
then have $\forall x \in ((\lambda v. v\$1) \text{'(path-image } p \cup \text{path-image } q)). ?b-x \geq x$
by (*simp add: assms(7) assms(8) bounded-component-cart bounded-has-Sup(1) bounded-simple-path-image*)
thus *?thesis using that[of ?b-x + 1] by (smt (verit, ccfv-SIG) assms(10) imageI)*
qed
obtain $b-y$ **where** $b-y: \forall v \in \text{path-image } p \cup \text{path-image } q. b-y > v\2
proof–
let $?b-y = \text{Sup } ((\lambda v. v\$2) \text{'(path-image } p \cup \text{path-image } q))$
have $\text{compact } (\text{path-image } p \cup \text{path-image } q)$
by (*simp add: assms(7) assms(8) compact-Un compact-simple-path-image*)
moreover have $\text{continuous-on UNIV } ((\lambda v. v\$2)::(\text{real}^2 \Rightarrow \text{real}))$
by (*simp add: continuous-on-component*)
ultimately have $*$: $\text{compact } ((\lambda v. v\$2) \text{'(path-image } p \cup \text{path-image } q))$
by (*meson compact-continuous-image continuous-on-subset top-greatest*)
then have $\forall x \in ((\lambda v. v\$2) \text{'(path-image } p \cup \text{path-image } q)). ?b-y \geq x$
by (*simp add: assms(7) assms(8) bounded-component-cart bounded-has-Sup(1) bounded-simple-path-image*)
thus *?thesis using that[of ?b-y + 1] by (smt (verit, ccfv-SIG) assms(10) imageI)*
qed

let $?l1 = \text{linepath } p1 \text{ (vector } [b-x, 0])$
let $?l2 = \text{linepath } (\text{vector } [b-x, 0]) \text{ ((vector } [b-x, b-y])::(\text{real}^2))$
let $?l3 = \text{linepath } (\text{vector } [b-x, b-y]) \text{ ((vector } [a-x, b-y])::(\text{real}^2))$
let $?l4 = \text{linepath } (\text{vector } [a-x, b-y]) \text{ ((vector } [a-x, 0])::(\text{real}^2))$
let $?l5 = \text{linepath } (\text{vector } [a-x, 0]) \text{ } p0$

let $?R' = ?l1 +++ ?l2 +++ ?l3 +++ ?l4 +++ ?l5$
let $?R = p +++ ?R'$

have $R-y-b: \forall v \in \text{path-image } ?R. v\$2 \leq b-y$
proof–
have $\forall v \in \text{path-image } ?l1. v\$2 \leq b-y$
by (*metis UnCI assms(9) b-y less-eq-real-def p1-def path-image-linepath pathfin-ish-in-path-image segment-horizontal vector-2(2)*)
moreover have $\forall v \in \text{path-image } ?l2. v\$2 \leq b-y$
by (*smt (verit, ccfv-SIG) UnCI assms(9) b-y p0-def path-image-linepath pathstart-in-path-image segment-vertical vector-2(1) vector-2(2)*)
moreover have $\forall v \in \text{path-image } ?l3. v\$2 \leq b-y$

by (*simp add: segment-horizontal*)
 moreover have $\forall v \in \text{path-image } ?l4. v \leq b-y$
 by (*smt (verit, best) UnCI assms(9) b-y p0-def path-image-linepath pathstart-in-path-image segment-vertical vector-2(1) vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l5. v \leq b-y$
 by (*smt (verit) UnI1 assms(9) b-y linepath-image-01 p0-def path-defs(4) pathstart-in-path-image segment-horizontal vector-2(2)*)
 ultimately show *?thesis* by (*smt (verit, best) UnCI b-y not-in-path-image-join*)
 qed
 have $R-y-q0: \forall v \in \text{path-image } ?R. v \geq q0$
 proof –
 have $\forall v \in \text{path-image } ?l1. v \geq q0$
 using *assms(13) assms(9) p1-def pathfinish-in-path-image segment-horizontal*
 by *fastforce*
 moreover have $\forall v \in \text{path-image } ?l2. v \geq q0$
 by (*smt (z3) UnCI assms(13) assms(9) b-y p1-def path-image-linepath pathfinish-in-path-image segment-vertical vector-2(1) vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l3. v \geq q0$
 by (*metis calculation(2) ends-in-segment(2) path-image-linepath segment-horizontal vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l4. v \geq q0$
 by (*smt (z3) UnCI assms(13) assms(9) b-y p1-def path-image-linepath pathfinish-in-path-image segment-vertical vector-2(1) vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l5. v \geq q0$
 by (*metis assms(13) assms(9) p0-def path-image-linepath pathstart-in-path-image segment-horizontal vector-2(2)*)
 ultimately show *?thesis*
 by (*metis assms(13) not-in-path-image-join*)
 qed

 have $R-x-a: \forall v \in \text{path-image } ?R. v \geq a-x$
 proof –
 have $\forall v \in \text{path-image } ?l1. v \geq a-x$
 by (*metis UnCI a-x assms(9) linorder-le-cases linorder-not-less p0-def path-image-linepath pathstart-in-path-image segment-horizontal vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l2. v \geq a-x$
 by (*smt (z3) UnCI assms(9) b-y calculation p0-def path-image-linepath pathstart-in-path-image pathstart-linepath segment-vertical vector-2(1) vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l3. v \geq a-x$
 by (*metis calculation(2) ends-in-segment(2) path-image-linepath segment-horizontal vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l4. v \geq a-x$
 by (*smt (z3) assms(9) calculation(1) calculation(3) ends-in-segment(1) path-image-linepath segment-vertical vector-2(1) vector-2(2)*)
 moreover have $\forall v \in \text{path-image } ?l5. v \geq a-x$
 by (*smt (verit, del-insts) UnCI a-x assms(9) p0-def path-image-linepath pathstart-in-path-image segment-horizontal vector-2(2)*)
 ultimately show *?thesis*
 by (*smt (z3) UnCI a-x assms(9) b-x not-in-path-image-join p1-def path-image-linepath*)

pathfinish-in-path-image segment-horizontal segment-vertical vector-2(1) vector-2(2)
qed

have *closed*: *closed-path ?R* **using** *assms p0-def unfolding simple-path-def closed-path-def*
by *simp*

have *simple*: *simple-path ?R*

proof–

have *arc ?R'*

proof–

let *?a = p1*

let *?b = (vector [b-x, 0])::(real^2)*

let *?c = (vector [b-x, b-y])::(real^2)*

let *?d = (vector [a-x, b-y])::(real^2)*

let *?e = (vector [a-x, 0])::(real^2)*

let *?f = p0*

have *arcs*: *arc ?l1 ∧ arc ?l2 ∧ arc ?l3 ∧ arc ?l4 ∧ arc ?l5*

by (*smt (verit, ccfv-SIG) UnCI a-x arc-linepath assms(9) b-x b-y p0-def*
p1-def pathfinish-in-path-image pathstart-in-path-image vector-2(1) vector-2(2))

have *l4l5*: *path-image ?l4 ∩ path-image ?l5 = {pathfinish ?l4}*

using *linepath-int-corner[of ?d ?e ?f] arc-simple-path arcs constant-linepath-is-not-loop-free*
p0 simple-path-def

by *auto*

have *l3l4*: *path-image ?l3 ∩ path-image ?l4 = {pathfinish ?l3}*

using *linepath-int-corner[of ?c ?d ?e]*

by (*metis Int-commute arc-simple-path arcs closed-segment-commute linepath-0'*
linepath-int-corner path-image-linepath pathfinish-linepath pathstart-def vector-2(2))

have *l2l3*: *path-image ?l2 ∩ path-image ?l3 = {pathfinish ?l2}*

using *linepath-int-corner[of ?b ?c ?d]*

by (*metis Int-commute arc-simple-path arcs linepath-0' linepath-int-corner*
pathfinish-linepath pathstart-def vector-2(2))

have *l1l2*: *path-image ?l1 ∩ path-image ?l2 = {pathfinish ?l1}*

using *linepath-int-corner[of ?a ?b ?c]*

by (*metis Int-commute arc-distinct-ends arcs assms(9) closed-segment-commute*
linepath-int-corner path-image-linepath pathfinish-linepath pathstart-linepath vector-2(2))

have *l3l5*: *path-image ?l3 ∩ path-image ?l5 = {}*

using *linepath-int-horizontal[of ?c ?d ?e ?f]*

by (*metis arc-distinct-ends arcs assms(9) linepath-int-horizontal pathfin-*
ish-linepath pathstart-linepath vector-2(2))

have *l2l4*: *path-image ?l2 ∩ path-image ?l4 = {}*

using *linepath-int-vertical[of ?b ?c ?d ?e]*

by (*metis arc-distinct-ends arcs linepath-int-vertical pathfinish-linepath path-*
start-linepath vector-2(1))

have *l1l3*: *path-image ?l1 ∩ path-image ?l3 = {}*

using *linepath-int-vertical[of ?a ?b ?c ?d]*

by (*metis arc-distinct-ends arcs assms(9) linepath-int-horizontal pathfin-*
ish-linepath pathstart-linepath vector-2(2))

have $l2l5$: $\text{path-image } ?l2 \cap \text{path-image } ?l5 = \{\}$
using $\text{linepath-int-columns}$ [of $?b ?c ?e ?f$]
by ($\text{smt (verit, ccfv-threshold) Int-commute UnCI a-x b-x linepath-int-columns}$
 $p0 p0\text{-def pathstart-in-path-image pathstart-join vector-2(1) verit-comp-simplify1(3)}$)
have $l1l4$: $\text{path-image } ?l1 \cap \text{path-image } ?l4 = \{\}$
using $\text{linepath-int-columns}$ [of $?a ?b ?d ?e$]
by ($\text{smt (z3) UnCI a-x assms(9) b-x disjoint-iff p1-def path-image-linepath}$
 $\text{pathfinish-in-path-image segment-horizontal segment-vertical vector-2(1) vector-2(2)}$)

have $l1l5$: $\text{path-image } ?l1 \cap \text{path-image } ?l5 = \{\}$
using $\text{linepath-int-columns}$ [of $?a ?b ?e ?f$]
by ($\text{smt (z3) UnCI a-def a-x assms(10) assms(9) b-x disjoint-iff p1-def}$
 $\text{path-image-linepath pathfinish-in-path-image segment-horizontal vector-2(1) vec}$
 tor-2(2))

have $\text{path-image } ?l4 \cap \text{path-image } ?l5 = \{\text{pathfinish } ?l4\}$
using $l4l5$ **by** blast
moreover **have** $\text{sf-45: pathfinish } ?l4 = \text{pathstart } ?l5$ **by** simp
ultimately **have** $\text{arc } (?l4 \text{ +++ } ?l5)$
by ($\text{metis arc-join-eq-alt arcs}$)
moreover **have** $\text{path-image } ?l3 \cap \text{path-image } (?l4 \text{ +++ } ?l5) = \{\text{pathfinish}$
 $?l3\}$
using $l3l4 l3l5$
by ($\text{metis (no-types, lifting) Int-Un-distrib sf-45 insert-is-Un path-image-join}$)
moreover **have** $\text{sf-345: pathfinish } ?l3 = \text{pathstart } (?l4 \text{ +++ } ?l5)$ **by** simp
ultimately **have** $\text{arc } (?l3 \text{ +++ } ?l4 \text{ +++ } ?l5)$
by ($\text{metis arc-join-eq-alt arcs}$)
moreover **have** $\text{path-image } ?l2 \cap \text{path-image } (?l3 \text{ +++ } ?l4 \text{ +++ } ?l5) =$
 $\{\text{pathfinish } ?l2\}$
using $l2l3 l2l4 l2l5$
by ($\text{smt (verit) Int-Un-distrib sf-45 sf-345 insert-is-Un path-image-join}$
 sup-bot-left)
moreover **have** $\text{sf-2345: pathfinish } ?l2 = \text{pathstart } (?l3 \text{ +++ } ?l4 \text{ +++ } ?l5)$
by simp
ultimately **have** $\text{arc } (?l2 \text{ +++ } ?l3 \text{ +++ } ?l4 \text{ +++ } ?l5)$
by ($\text{metis arc-join-eq-alt arcs}$)
moreover **have** $\text{path-image } ?l1 \cap \text{path-image } (?l2 \text{ +++ } ?l3 \text{ +++ } ?l4 \text{ +++}$
 $?l5) = \{\text{pathfinish } ?l1\}$
proof–
have $\text{path-image } (?l2 \text{ +++ } ?l3 \text{ +++ } ?l4 \text{ +++ } ?l5)$
 $= \text{path-image } ?l2 \cup \text{path-image } ?l3 \cup \text{path-image } ?l4 \cup \text{path-image } ?l5$
by ($\text{simp add: path-image-join sup-assoc}$)
thus $?thesis$ **using** $l1l2 l1l3 l1l4 l1l5$ **by** blast
qed
moreover **have** $\text{pathfinish } ?l1 = \text{pathstart } (?l2 \text{ +++ } ?l3 \text{ +++ } ?l4 \text{ +++}$
 $?l5)$ **by** simp
ultimately **show** $\text{arc } (?l1 \text{ +++ } ?l2 \text{ +++ } ?l3 \text{ +++ } ?l4 \text{ +++ } ?l5)$
by ($\text{metis arc-join-eq-alt arcs}$)

qed
moreover have *loop-free p* **using** *assms(1) assms(7) simple-path-def* **by** *blast*
moreover have *path-image ?R' ∩ path-image p = {p0, p1}*
proof-
have *path-image p ∩ path-image ?l2 = {}* **using** *b-x segment-vertical* **by** *auto*
moreover have *path-image p ∩ path-image ?l3 = {}* **using** *b-y segment-horizontal*
by *auto*
moreover have *path-image p ∩ path-image ?l4 = {}* **using** *a-x segment-vertical*
by *auto*
moreover have *path-image p ∩ path-image ?l1 = {p1}*
proof-
have *p1 ∈ path-image p* **using** *p1-def* **by** *blast*
moreover have *path-image p ∩ path-image ?l1 ⊆ {p1}*
proof(*rule subsetI*)
fix *x* **assume** ***: *x ∈ path-image p ∩ path-image ?l1*
then have *x\$1 ≤ a*
using *a-def assms(10) assms(12) assms(9) l-def linepath-image-01*
segment-horizontal **by** *auto*
moreover have *x\$1 ≥ a*
by (*smt (z3) * Int-iff Un-iff a-def assms(9) b-x linepath-image-01*
path-defs(4) segment-horizontal vector-2(1) vector-2(2))
moreover have *x\$2 = 0* **using** ** assms(9) segment-horizontal* **by** *auto*
ultimately show *x ∈ {p1}* **using** *a-def assms(9) segment-vertical* **by**
fastforce
qed
ultimately show *?thesis* **by** *auto*
qed
moreover have *path-image p ∩ path-image ?l5 = {p0}*
proof-
have *p0 ∈ path-image p* **using** *p0-def* **by** *blast*
moreover have *path-image p ∩ path-image ?l5 ⊆ {p0}*
proof(*rule subsetI*)
fix *x* **assume** ***: *x ∈ path-image p ∩ path-image ?l5*
then have *x\$1 ≤ 0*
using *R-x-a assms(9) p0-def pathstart-in-path-image segment-horizontal*
by *fastforce*
moreover have *x\$1 ≥ 0*
proof-
have *x ∈ {x. x\$2 = 0}* **using** ** assms(9) segment-horizontal* **by** *fastforce*
then have *x ∈ l* **using** ** assms(12)* **by** *auto*
thus *?thesis* **using** *a-def assms(10) assms(9) l-def segment-horizontal*
by *auto*
qed
moreover have *x\$2 = 0* **using** ** assms(9) segment-horizontal* **by** *auto*
ultimately show *x ∈ {p0}* **using** *a-def assms(9) segment-vertical* **by**
fastforce
qed
ultimately show *?thesis* **by** *auto*
qed

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moreover have path-image ?R'
  = path-image ?l1 ∪ path-image ?l2 ∪ path-image ?l3 ∪ path-image ?l4 ∪
  path-image ?l5
  by (simp add: Un-assoc path-image-join)
  ultimately show ?thesis by fast
qed
moreover have arc p
  using a-def arc-simple-path assms(10) assms(7) p0 p0-def p1-def by fastforce
  ultimately show ?thesis
  by (metis (no-types, lifting) simple-path-join-loop-eq Int-commute dual-order.refl
  p0-def p1-def pathfinish-join pathfinish-linepath pathstart-join pathstart-linepath)
qed

have inside-outside: inside-outside ?R (path-inside ?R) (path-outside ?R)
  using closed simple Jordan-inside-outside-real2
  by (simp add: closed-path-def inside-outside-def path-inside-def path-outside-def)

have interior-frontier: path-inside ?R = interior (path-inside ?R)
  ∧ frontier (path-inside ?R) = path-image ?R
  using inside-outside interior-open unfolding inside-outside-def by auto

have path-image q ∩ path-image ?l1 ⊆ {p1}
proof(rule subsetI)
  fix x assume *: x ∈ path-image q ∩ path-image ?l1
  then have x$1 ≤ a using a-def assms(10) assms(11) assms(9) l-def seg-
  ment-horizontal by auto
  moreover have x$1 ≥ a
  by (smt (z3) * Int-iff Un-iff a-def assms(9) b-x linepath-image-01 path-defs(4)
  segment-horizontal vector-2(1) vector-2(2))
  moreover have x$2 = 0 using * assms(9) segment-horizontal by auto
  ultimately show x ∈ {p1} using a-def assms(9) segment-vertical by fastforce
qed
moreover have path-image q ∩ path-image ?l5 ⊆ {p0}
proof(rule subsetI)
  fix x assume *: x ∈ path-image q ∩ path-image ?l5
  then have x$1 ≤ 0
  using R-x-a assms(9) p0-def pathstart-in-path-image segment-horizontal by
  fastforce
  moreover have x$1 ≥ 0
  using * a-def assms(10) assms(11) assms(9) l-def segment-horizontal by auto
  moreover have x$2 = 0 using * assms(9) segment-horizontal by auto
  ultimately show x ∈ {p0} using a-def assms(9) segment-vertical by fastforce
qed
moreover have ?thesis if p1 ∈ path-image q ∩ path-image ?l1 using p1-def that
by blast
moreover have ?thesis if p0 ∈ path-image q ∩ path-image ?l5 using p0-def that
by blast
moreover have ?thesis if
  q-int-l1: path-image q ∩ path-image ?l1 = {} and

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    q-int-l5: path-image q ∩ path-image ?l5 = {}
  proof-
    have q-int-l2: path-image q ∩ path-image ?l2 = {}
      using b-x segment-vertical by auto
    moreover have q-int-l3: path-image q ∩ path-image ?l3 = {}
      using UnCI b-y segment-horizontal by auto
    moreover have q-int-l4: path-image q ∩ path-image ?l4 = {}
      using a-x segment-vertical by auto
    moreover have ?thesis if q0 ∈ path-image p using q0-def that by blast
    moreover have path-image q ∩ path-image ?R ≠ {} if q0 ∉ path-image p
  proof-
    have q0 ∈ path-outside ?R

  proof-
    let ?e2' = (vector [0, -1])::(real^2)
    let ?ray = λd. q0 + d *R ?e2'
    have ¬ (∃ d>0. ?ray d ∈ path-image ?R)
  proof-
    have ∀ d>0. (?ray d)$2 < q0$2 by auto
    thus ?thesis using R-y-q0 by fastforce
  qed
  moreover have bounded (path-inside ?R) using bounded-finite-inside simple
by blast
  moreover have ?e2' ≠ 0 by (metis vector-2(2) zero-index zero-neg-neg-one)
  ultimately have q0 ∉ path-inside ?R
    using ray-to-frontier[of path-inside ?R] interior-frontier by metis
  moreover have q0 ∉ path-image ?R
    using that q-int-l1 q-int-l2 q-int-l3 q-int-l4 q-int-l5
    by (simp add: disjoint-iff not-in-path-image-join pathstart-in-path-image
q0-def)
  ultimately show ?thesis using inside-outside unfolding inside-outside-def
by blast
  qed
  then have q0 ∈ - (path-inside ?R)
by (metis ComplI IntI equals0D inside-Int-outside path-inside-def path-outside-def)
  moreover have q1 ∈ path-inside ?R

  proof-
    let ?e = (vector [q1$1, b-y])::(real^2)
    let ?d1 = (vector [b-x, b-y])::(real^2)
    let ?d2 = (vector [a-x, b-y])::(real^2)
    obtain ε where ε: 0 < ε ∧ ε < dist ?e q1 ∧ ε < dist ?e ?d1 ∧ ε < dist ?e
?d2
  proof-
    have ?e ≠ q1
      by (metis UnCI b-y order-less-irrefl pathfinish-in-path-image q1-def
vector-2(2))
    moreover have ?e ≠ ?d1
      by (smt (verit) UnCI b-x pathfinish-in-path-image q1-def vector-2(1))

```

moreover have $?e \neq ?d2$
by (*metis UnCI a-x order-less-irrefl pathfinish-in-path-image q1-def vector-2(1)*)
ultimately have $0 < \text{dist } ?e \ q1 \wedge 0 < \text{dist } ?e \ ?d1 \wedge 0 < \text{dist } ?e \ ?d2$ **by**
simp
then have $0 < \text{Min } \{\text{dist } ?e \ q1, \text{dist } ?e \ ?d1, \text{dist } ?e \ ?d2\}$ **by** *auto*
then obtain ε **where** $0 < \varepsilon \wedge \varepsilon < \text{Min } \{\text{dist } ?e \ q1, \text{dist } ?e \ ?d1, \text{dist } ?e \ ?d2\}$
by (*meson field-lbound-gt-zero*)
thus *?thesis* **using that** **by** *auto*
qed
then have $?e \in \text{path-image } ?l3$
by (*simp add: a-x b-x q1-def segment-horizontal less-eq-real-def pathfinish-in-path-image*)
then have $?e \in \text{path-image } ?R$ **by** (*simp add: p1-def path-image-join*)
then have $?e \in \text{frontier } (\text{path-inside } ?R)$
using *inside-outside unfolding inside-outside-def* **by** *blast*
then obtain *int-p* **where** $\text{int-p} : \text{int-p} \in \text{ball } ?e \ \varepsilon \wedge \text{int-p} \in \text{path-inside } ?R$
by (*meson \varepsilon inside-outside frontier-straddle mem-ball*)

have $\text{int-p-x} : a-x < \text{int-p}\$1 \wedge \text{int-p}\$1 < b-x$
by (*metis (mono-tags, lifting) dist-bound-1 UnI2 \varepsilon a-x b-x dist-commute int-p pathfinish-in-path-image q1-def vector-2(1) vector-2(2)*)
have $\text{int-p}\$2 < b-y$
proof(*rule ccontr*)
have $\text{int-p}\$2 \neq b-y$
proof–
have $\text{int-p}\$2 = b-y \implies \text{int-p} \in \text{path-image } ?l3$
using *int-p-x* **by** (*simp add: segment-horizontal*)
moreover have $\text{int-p} \in \text{path-image } ?l3 \implies \text{int-p} \in \text{path-image } ?R$
by (*simp add: p1-def path-image-join*)
moreover have $\text{path-image } ?R \cap \text{path-inside } ?R = \{\}$
using *inside-outside unfolding inside-outside-def* **by** *blast*
ultimately show *?thesis* **using** *int-p* **by** *fast*
qed
moreover assume $\neg \text{int-p}\$2 < b-y$
ultimately have $*$: $\text{int-p}\$2 > b-y$ **by** *simp*

let $?e2 = (\text{vector } [0, 1]) :: (\text{real}^2)$
let $?ray = \lambda d. \text{int-p} + d *_{\mathbb{R}} ?e2$
have $\neg (\exists d > 0. ?ray \ d \in \text{path-image } ?R)$
proof–
have $\forall d > 0. (?ray \ d)\$2 > b-y$ **using** $*$ **by** *auto*
thus *?thesis* **using** *R-y-b* **by** *fastforce*
qed
moreover have *bounded* (*path-inside* $?R$) **using** *bounded-finite-inside simple* **by** *blast*
moreover have $?e2 \neq 0$ **using** *e1e2-basis(4)* **by** *force*
ultimately have $\text{int-p} \notin \text{path-inside } ?R$

```

    using ray-to-frontier[of path-inside ?R] interior-frontier by metis
    thus False using int-p by blast
qed
moreover have int-p$2 > q1$2
proof-
  have dist int-p ?e < ε using ε dist-commute-lessI int-p mem-ball by blast
  then have dist (int-p$2) (?e$2) < ε by (smt (verit, best) dist-vec-nth-le)
  then have 1: int-p$2 > ?e$2 - ε by (simp add: dist-real-def)

  have q1$1 = ?e$1 by simp
  then have dist q1 ?e = dist (q1$2) (?e$2) using axis-dist by blast
  then have q1$2 < ?e$2 - ε
  by (smt (verit) UnCI ε b-y dist-commute dist-real-def pathfinish-in-path-image
q1-def vector-2(2))
  moreover have q1$2 < ?e$2 by (simp add: b-y pathfinish-in-path-image
q1-def)
  moreover have dist q1 ?e > ε by (metis ε dist-commute)
  ultimately have q1$2 < ?e$2 - ε by presburger
  thus ?thesis using 1 by force
qed
ultimately have int-p-y: int-p$2 < b-y ∧ int-p$2 > q1$2 by blast

let ?int-l = linepath int-p q1

have path-image ?int-l ∩ path-image p = {}
proof-
  have ∀ x ∈ path-image p. (?int-l 0)$2 > x$2
  by (smt (verit) int-p-y assms(14) linepath-0')
  moreover have ∀ x ∈ path-image p. (?int-l 1)$2 > x$2
  by (simp add: assms(14) linepath-1')
  ultimately have ∀ x ∈ path-image p. ∀ y ∈ path-image ?int-l. y$2 > x$2
  by (metis assms(14) linepath-0' linepath-bound-2(1))
  thus ?thesis by blast
qed
moreover have path-image ?int-l ∩ path-image ?l1 = {}
by (smt (verit, best) assms(14) assms(9) disjoint-iff int-p-y linepath-int-rows
p0-def pathstart-in-path-image vector-2(2))
moreover have path-image ?int-l ∩ path-image ?l2 = {}
by (metis UnCI b-x int-p-x linepath-int-columns pathfinish-in-path-image
q1-def vector-2(1))
moreover have path-image ?int-l ∩ path-image ?l3 = {}
using int-p-y linepath-int-rows by auto
moreover have path-image ?int-l ∩ path-image ?l4 = {}
by (metis UnCI a-x inf-commute int-p-x linepath-int-columns pathfin-
ish-in-path-image q1-def vector-2(1))
moreover have path-image ?int-l ∩ path-image ?l5 = {}
by (smt (verit, best) assms(14) assms(9) disjoint-iff int-p-y linepath-int-rows
p0-def pathstart-in-path-image vector-2(2))
ultimately have path-image ?int-l ∩ path-image ?R = {}

```

by (*simp add: disjoint-iff not-in-path-image-join*)
then have $\text{path-image } ?\text{int-}l \subseteq \text{path-inside } ?R \vee \text{path-image } ?\text{int-}l \subseteq$
 $\text{path-outside } ?R$
by (*smt (verit, ccfv-SIG) convex-imp-path-connected convex-segment(1) disjoint-insert(1) insert-Diff inside-outside-def int-p linepath-image-01 local.inside-outside path-connected-not-frontier-subset path-defs(4) pathstart-in-path-image pathstart-linepath*)
moreover have $?\text{int-}l \cap \text{int-}p \in \text{path-inside } ?R$
using *int-p* **by** (*simp add: linepath-0'*)
ultimately have $\text{path-image } ?\text{int-}l \subseteq \text{path-inside } ?R$
using *inside-outside-def local.inside-outside* **by auto**
thus *?thesis* **by auto**
qed
ultimately have $\text{path-image } q \cap -(\text{path-inside } ?R) \neq \{\} \wedge \text{path-image } q \cap$
 $(\text{path-inside } ?R) \neq \{\}$
unfolding *q0-def q1-def* **by fast**
moreover have *path-connected (path-image q)*
by (*simp add: assms(8) path-connected-path-image simple-path-imp-path*)
moreover have $\text{path-image } ?R = \text{frontier } (\text{path-inside } ?R)$
using *inside-outside unfolding inside-outside-def p0-def path-inside-def* **by auto**
ultimately show *?thesis* **by** (*metis Diff-eq Diff-eq-empty-iff path-connected-not-frontier-subset*)
qed
ultimately show *?thesis*
by (*smt (verit, ccfv-threshold) disjoint-iff-not-equal not-in-path-image-join q-int-l1 q-int-l5*)
qed
ultimately show *?thesis* **by auto**
qed

lemma *pocket-fill-line-int-aux2:*

fixes $p \ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)\ \text{set}$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $a \equiv p1\$1$
defines $l \equiv \text{closed-segment } p0 \ p1$
assumes *simple-path p*
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes *convex A* \wedge *compact A*
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p \ \{0 <..<1\} \subseteq \text{interior } A$
shows $\text{path-image } p \cap \{x. x\$2 = 0\} \subseteq l$

proof –

have $l = \{x. x\$2 = 0 \wedge x\$1 \in \{0..a\}\}$
using *horizontal-segment-at-0' a-def assms(6) assms(7) l-def* **by presburger**
have *endpoints: (p 0)\$1 = 0* \wedge *(p 0)\$2 = 0* \wedge *(p 1)\$1 = a* \wedge *(p 1)\$2 = 0*
by (*metis a-def assms(6) p0-def p1-def pathfinish-def pathstart-def*)

have *False* **if** $*$: $\exists t \in \{0..1\}. (p\ t)\$2 = 0 \wedge ((p\ t)\$1 > a \vee (p\ t)\$1 < 0)$
proof–
obtain t **where** $t \in \{0 < .. < 1\} \wedge (p\ t)\$2 = 0 \wedge ((p\ t)\$1 > a \vee (p\ t)\$1 < 0)$
by (*metis * assms(7) endpoints atLeastAtMost-iff greaterThanLessThan-iff less-eq-real-def linorder-not-le*)
then obtain x **where** $x: x \in p'\{0 < .. < 1\} \wedge x\$2 = 0 \wedge (x\$1 > a \vee x\$1 < 0)$
by *blast*
thus *False*
using *pocket-fill-line-int-aux*[of $p0\ p1\ x\ A$]
by (*smt (verit, del-insts) Diff-iff a-def assms(10) assms(6) assms(7) assms(8) assms(9) empty-subsetI endpoints exhaust-2 frontier-def frontier-subset-compact insert-subset interior-subset p0-def pathstart-def subset-eq vec-eq-iff zero-index*)
qed
then have $\forall t \in \{0..1\}. (p\ t)\$2 = 0 \longrightarrow (p\ t)\$1 \in \{0..a\}$ **by** *fastforce*
then have $\forall v \in \text{path-image } p. v\$2 = 0 \longrightarrow v\$1 \in \{0..a\}$ **by** (*simp add: imageE path-defs(4)*)
thus *?thesis* **using** l **by** *blast*
qed

lemma *three-points-on-line*:

fixes $a\ b :: 'a::\text{real-vector}$
assumes $A = \text{affine hull } \{a, b\}$
assumes $a \neq b$
assumes $\{x, y, z\} \subseteq A$
assumes $x \neq y \wedge y \neq z \wedge x \neq z$
shows $x \in \text{open-segment } y\ z \vee y \in \text{open-segment } x\ z \vee z \in \text{open-segment } x\ y$
proof–
let $?u = b - a$

have $*$: $\bigwedge \alpha\ \beta\ \gamma :: \text{real}. \alpha \in \text{open-segment } \beta\ \gamma$
 $\implies a + \alpha *_R ?u \in \text{open-segment } (a + \beta *_R ?u)\ (a + \gamma *_R ?u)$

proof–
fix $\alpha\ \beta\ \gamma :: \text{real}$
assume $*$: $\alpha \in \text{open-segment } \beta\ \gamma$

define x **where** $x \equiv a + \alpha *_R ?u$
define y **where** $y \equiv a + \beta *_R ?u$
define z **where** $z \equiv a + \gamma *_R ?u$

obtain v **where** $v: \alpha = (1 - v) * \beta + v * \gamma \wedge v \in \{0 < .. < 1\}$
by (*metis (no-types, lifting) * imageE in-segment(2) real-scaleR-def segment-image-interval(2)*)
then have $x = a + ((1 - v) * \beta + v * \gamma) *_R ?u$ **using** x -**def** **by** *blast*
also have $\dots = a + (((1 - v) * \beta) *_R ?u) + ((v * \gamma) *_R ?u)$ **by** (*simp add: scaleR-left.add*)
also have $\dots = a + ((1 - v) *_R (\beta *_R ?u)) + (v *_R (\gamma *_R ?u))$ **by** *simp*
also have $\dots = a + ((1 - v) *_R (y - a)) + (v *_R (z - a))$ **by** (*simp add: y-def z-def*)
also have $\dots = a + y - a - v *_R (y - a) + v *_R (z - a)$ **by** (*simp add:*

scaleR-left-diff-distrib
also have $\dots = y - v *_R (y - a) + v *_R (z - a)$ **by** *simp*
also have $\dots = y - (v *_R y) + (v *_R a) + (v *_R z) - (v *_R a)$ **by** (*simp add: scaleR-right-diff-distrib*)
also have $\dots = (1 - v) *_R y + v *_R z$ **by** (*metis add-diff-cancel diff-add-eq scaleR-collapse*)
finally have $x = (1 - v) *_R y + v *_R z$.
moreover have $0 \leq 1 - v \wedge 1 - v \leq 1$ **using** v **by** *fastforce*
ultimately have $x \in \text{closed-segment } y z$ **using** *in-segment(1)* **by** *auto*
moreover have $x \neq y \wedge x \neq z$
by (*metis * add-diff-cancel-left' assms(2) eq-iff-diff-eq-0 in-open-segment-iff-line open-segment-commute open-segment-subsegment scaleR-right-imp-eq x-def y-def z-def*)
ultimately show $a + \alpha *_R ?u \in \text{open-segment } (a + \beta *_R ?u) (a + \gamma *_R ?u)$
unfolding *open-segment-def* **using** *x-def y-def z-def* **by** *force*
qed

obtain $\alpha \beta \gamma$ **where** $xyz: x = a + \alpha *_R ?u \wedge y = a + \beta *_R ?u \wedge z = a + \gamma *_R ?u$
using *affine-hull-2-alt[of a b] assms(1) assms(3)* **by** *auto*
then have $\alpha \neq \beta \wedge \beta \neq \gamma \wedge \alpha \neq \gamma$ **using** *assms* **by** *blast*
moreover have $\alpha \in \text{closed-segment } \beta \gamma \vee \beta \in \text{closed-segment } \alpha \gamma \vee \gamma \in \text{closed-segment } \alpha \beta$
by (*metis atLeastAtMost-iff closed-segment-commute less-eq-real-def less-max-iff-disj linorder-not-less real-Icc-closed-segment*)
ultimately have $\alpha \in \text{open-segment } \beta \gamma \vee \beta \in \text{open-segment } \alpha \gamma \vee \gamma \in \text{open-segment } \alpha \beta$
unfolding *open-segment-def* **by** *fast*
thus *?thesis* **using** ** xyz* **by** *presburger*
qed

lemma *pocket-fill-line-int-aux3*:

fixes $A :: (\text{real}^2)$ *set*
assumes *convex A* \wedge *compact A*
assumes $v \neq 0$
assumes *closed-segment* $0 w \subseteq \text{frontier } A$ (**is** *closed-segment* $?a ?b \subseteq -$)
assumes $w \cdot v = 0$
assumes $w \neq 0$
shows $(A \subseteq \{x. x \cdot v \leq 0\} \vee A \subseteq \{x. x \cdot v \geq 0\})$ (**is** $A \subseteq ?P1 \vee A \subseteq ?P2$)
proof –
have *frontiers*: $\text{frontier } ?P1 = \text{frontier } ?P2 \wedge \text{frontier } ?P1 \subseteq ?P2 \wedge \text{frontier } ?P2 \subseteq ?P1$
by (*smt (verit, ccfv-threshold) Collect-mono assms(2) frontier-halfspace-component-ge frontier-halfspace-le inner-commute subset-antisym*)
have *frontier*: $\text{frontier } ?P1 = \{x. x \cdot v = 0\}$
by (*simp add: assms(2) frontier-halfspace-component-ge frontiers*)

have *?thesis* **if** $\text{interior } A \neq \{\}$
proof –
have $\text{interior } A \subseteq ?P1 \vee \text{interior } A \subseteq ?P2$

proof(*rule ccontr*)
assume \neg (*interior* $A \subseteq ?P1 \vee$ *interior* $A \subseteq ?P2$)
then obtain $x y$ **where** $xy: x \in ((\text{interior } A) \cap ?P1) - ?P2 \wedge y \in ((\text{interior } A) \cap ?P2) - ?P1$
by *fastforce*
moreover have $x \in \text{frontier } ?P1 \cup \text{interior } ?P1 \wedge y \in \text{frontier } ?P2 \cup \text{interior } ?P2$
by (*metis DiffD1 IntD2 Un-Diff-cancel2 frontiers closure-Un-frontier frontier-def interior-subset sup.orderE xy*)
ultimately have $xy': x \in (\text{interior } A) \cap \text{interior } ?P1 \wedge y \in (\text{interior } A) \cap \text{interior } ?P2$
using *frontiers by blast*
then have *closed-segment* $x y \cap \text{frontier } ?P1 \neq \{\}$
by (*metis (no-types, lifting) DiffD1 DiffD2 Int-iff convex-closed-segment convex-imp-path-connected empty-iff ends-in-segment(1) ends-in-segment(2) in-mono path-connected-not-frontier-subset xy*)
moreover have *closed-segment* $x y \subseteq \text{interior } A$
by (*metis convex-interior Int-iff assms(1) convex-contains-segment xy*)
ultimately obtain z **where** $z: z \in \text{interior } A \cap \text{frontier } ?P1$ **by** *blast*

have *closed-segment* $?a ?b \subseteq \text{frontier } ?P1$
proof(*rule subsetI*)
fix x
assume $x \in \text{closed-segment } ?a ?b$
then obtain u **where** $x = (1 - u) *_R ?a + u *_R ?b \wedge 0 \leq u \wedge u \leq 1$
unfolding *closed-segment-def* **by** *blast*
then have $x \cdot v = u *_R (?b \cdot v)$ **by** *simp*
moreover have $?b \cdot v = 0$ **by** (*simp add: assms(4)*)
ultimately have $x \cdot v = 0$ **by** *simp*
thus $x \in \text{frontier } ?P1$ **using** *frontier* **by** *blast*

qed
moreover have $z \notin \text{closed-segment } ?a ?b$ **using** *assms(3) frontier-def z* **by** *fastforce*
ultimately have $z \in \text{frontier } ?P1 - \text{closed-segment } ?a ?b$ **using** z **by** *blast*
moreover have *collinear* $\{z, ?a, ?b\}$
proof–
have $\{z, ?a, ?b\} \subseteq \{x. x \cdot v = 0\}$
using $\langle \{0--w\} \subseteq \text{frontier } \{x. x \cdot v \leq 0\} \rangle$ *frontier z* **by** *auto*
moreover have $\{x. x \cdot v = 0\} = \text{affine hull } \{?a, ?b\}$
by (*metis (no-types, lifting) Collect-mono assms(2) assms(5) calculation halfplane-frontier-affine-hull inner-commute insert-subset subset-antisym*)
ultimately show *thesis* **using** *collinear-affine-hull* **by** *auto*

qed
ultimately have $?a \in \text{open-segment } z ?b \vee ?b \in \text{open-segment } z ?a$
using *three-points-on-line*[of $\{x. x \cdot v = 0\}$]
by (*smt (z3) $\langle z \notin \{0--w\} \rangle$ *assms(5) collinear-3-imp-in-affine-hull ends-in-segment(1) ends-in-segment(2) hull-redundant hull-subset insert-commute open-closed-segment three-points-on-line*)*
moreover have *open-segment* $z ?b \subseteq \text{interior } A \wedge$ *open-segment* $z ?a \subseteq$

interior A

proof–

have closed-segment z ?b \subseteq A \wedge closed-segment z ?a \subseteq A

by (meson IntD1 assms(1) assms(3) closed-segment-subset ends-in-segment(1) ends-in-segment(2) frontier-subset-compact in-mono interior-subset z)

then have rel-interior (closed-segment z ?b) \subseteq interior A

\wedge rel-interior (closed-segment z ?a) \subseteq interior A

by (metis IntD1 $\langle z \notin \{0--w\} \rangle$ assms(1) closure-convex-hull convex-hull-eq in-interior-closure-convex-segment order-class.order-eq-iff rel-interior-closed-segment subsetD subset-closed-segment z)

moreover have rel-interior (closed-segment z ?b) = open-segment z ?b

\wedge rel-interior (closed-segment z ?a) = open-segment z ?a

by (metis $\langle z \notin \{0--w\} \rangle$ closed-segment-commute ends-in-segment(1) rel-interior-closed-segment)

ultimately show ?thesis **by force**

qed

ultimately have ?a \in interior A \vee ?b \in interior A **by fast**

thus False **using** assms(3) frontier-def **by auto**

qed

then have closure (interior A) \subseteq closure ?P1 \vee closure (interior A) \subseteq closure ?P2

using closure-mono **by blast**

moreover have closed ?P1 \wedge closed ?P2

by (simp add: closed-halfspace-component-ge closed-halfspace-component-le)

moreover have closure (interior A) = A

using assms(1)

by (simp add: compact-imp-closed convex-closure-interior that)

ultimately show ?thesis **using** closure-closed **by auto**

qed

moreover have ?thesis **if** interior A = {}

proof(rule ccontr)

assume \neg (A \subseteq ?P1 \vee A \subseteq ?P2)

then obtain x y **where** xy: x \in (A \cap ?P1) $-$?P2 \wedge y \in (A \cap ?P2) $-$?P1

by fastforce

moreover have x \in frontier ?P1 \cup interior ?P1 \wedge y \in frontier ?P2 \cup interior ?P2

by (metis DiffD1 IntD2 Un-Diff-cancel2 frontiers closure-Un-frontier frontier-def interior-subset sup.orderE xy)

ultimately have xy': x \in A \cap interior ?P1 \wedge y \in A \cap interior ?P2 **using** frontiers **by blast**

have \neg collinear {?a, ?b, x, y}

proof(rule ccontr)

assume $\neg \neg$ collinear {?a, ?b, x, y}

then have *: collinear {?a, ?b, x, y} **by blast**

then have {?a, ?b, x, y} \subseteq affine hull {?a, ?b}

by (metis assms(5) collinear-3-imp-in-affine-hull collinear-4-3 hull-subset insert-subset)

moreover have affine hull {?a, ?b} = {x. x \cdot v = 0}

by (smt (verit) DiffE * assms(2) assms(4) assms(5) collinear-3-imp-in-affine-hull

collinear-4-3 halfplane-frontier-affine-hull inner-commute mem-Collect-eq xy
moreover have ... = frontier ?P1 \wedge ... = frontier ?P2
using frontiers *assms(2) frontier-halfspace-component-ge* **by** blast
ultimately show False **using** frontiers *xy* **by** auto
qed
then obtain *c1 c2 c3* **where** *c123*: \neg collinear {*c1, c2, c3*} \wedge {*c1, c2, c3*}
 \subseteq {*?a, ?b, x, y*}
by (*metis assms(5) collinear-4-3 insert-mono subset-insertI*)
then have interior (convex hull {*c1, c2, c3*}) \neq {}
by (*metis Jordan-inside-outside-real2 closed-path-def make-triangle-def path-inside-def*
polygon-def polygon-of-def triangle-inside-is-convex-hull-interior triangle-is-polygon)
moreover have {*c1, c2, c3*} \subseteq *A*
by (*smt (verit, del-insts) c123 xy' assms(1) assms(3) empty-subsetI fron-*
tier-subset-compact in-mono inf.orderE insert-absorb insert-mono le-infE subsetI
subset-closed-segment)
ultimately have interior *A* \neq {}
by (*metis assms(1) interior-mono subset-empty subset-hull*)
thus False **using** that **by** blast
qed
ultimately show *?thesis* **by** blast
qed

lemma *pocket-fill-line-int-aux4*:

fixes *p q* :: *R-to-R2*
fixes *A* :: (*real*²) set
defines *p0* \equiv pathstart *p*
defines *p1* \equiv pathfinish *p*
defines *q0* \equiv pathstart *q*
defines *q1* \equiv pathfinish *q*
defines *a* \equiv *p1*\$1
defines *l* \equiv closed-segment *p0 p1*
assumes simple-path *p*
assumes simple-path *q*
assumes path-image *p* \cap path-image *q* = {}
assumes *p0*\$1 = 0 \wedge *p0*\$2 = 0 \wedge *p1*\$2 = 0
assumes *a* > 0
assumes $\forall v \in$ path-image *p*. *q0*\$2 \leq *v*\$2
assumes $\forall v \in$ path-image *p*. *q1*\$2 > *v*\$2
assumes convex *A* \wedge compact *A*
assumes {*p0, p1*} \subseteq frontier *A*
assumes *p*{0<..*1*} \subseteq interior *A*
assumes path-image *q* \subseteq *A*
shows *l* \subseteq frontier *A* $\forall x \in$ (path-image *p*) \cup (path-image *q*). *x*\$2 \geq 0 *q0*\$2 = 0
proof –
have *l* = {*x*. *x*\$2 = 0 \wedge *x*\$1 \in {0..*a*}}
using horizontal-segment-at-0' *a-def assms(10) assms(11) l-def* **by** presburger
have endpoints: (*p 0*)\$1 = 0 \wedge (*p 0*)\$2 = 0 \wedge (*p 1*)\$1 = *a* \wedge (*p 1*)\$2 = 0
by (*metis a-def assms(10) p0-def p1-def pathfinish-def pathstart-def*)

have $l \subseteq \text{frontier } A$ **if** $\neg (\text{path-image } q \cap \{x. x\$2 = 0\} \subseteq l)$
proof –
from that obtain x **where** $x \in \text{path-image } q \cap \{x. x\$2 = 0\} \wedge (x\$1 < 0 \vee x\$1 > a)$
by (*smt (verit) Int-Collect a-def assms(10) endpoints l-def p0-def pathstart-def segment-horizontal subsetI*)
thus *?thesis*
using *pocket-fill-line-int-aux*[of $p0$ $p1$ x A] **unfolding** *l-def*
by (*smt (verit, del-insts) IntD2 Int-commute a-def assms(11) assms(14) assms(15) assms(17) assms(10) endpoints exhaust-2 frontier-subset-compact insert-subset mem-Collect-eq p0-def pathstart-def subset-eq vec-eq-iff zero-index*)
qed
moreover have *False* **if** $(\text{path-image } q \cap \{x. x\$2 = 0\} \subseteq l)$
proof –
have $(\text{path-image } p \cap \{x. x\$2 = 0\} \subseteq l)$
using *pocket-fill-line-int-aux2*
by (*metis a-def assms(10) assms(11) assms(14) assms(15) assms(16) assms(7) l-def p0-def p1-def*)
then have $\text{path-image } p \cap \text{path-image } q \neq \{\}$
using *pocket-fill-line-int-aux1*
by (*metis (mono-tags, lifting) assms(11) assms(12) assms(13) assms(7) assms(8) endpoints l-def p0-def p1-def pathfinish-def pathstart-def q0-def q1-def that*)
thus *False* **by** (*simp add: assms(9)*)
qed
ultimately show $*$: $l \subseteq \text{frontier } A$ **by** *blast*

show $\forall x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 \geq 0$
proof(*rule ccontr*)
assume $\neg (\forall x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 \geq 0)$
then have $\exists x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 < 0$ **using** *linorder-not-le*
by *blast*
then obtain x **where** $x: x \in ((\text{path-image } p) \cup (\text{path-image } q)) \cap A \wedge x\$2 < 0$
using *assms(12) assms(17) pathstart-in-path-image q0-def* **by** *fastforce*

let $?v = (\text{vector } [0, 1])::(\text{real}^2)$
have $1: ?v \neq 0$ **by** (*simp add: e1e2-basis(3)*)
have $2: \text{closed-segment } 0 \ p1 \subseteq \text{frontier } A$
by (*smt (verit, del-insts) * Int-closed-segment closed-segment-eq doubleton-eq-iff endpoints l-def p0-def pathstart-def segment-vertical zero-index*)
have $3: p1 \cdot ?v = 0$ **by** (*metis assms(10) cart-eq-inner-axis e1e2-basis(3)*)
have $4: p1 \neq 0$ **using** *a-def assms(11)* **by** *force*
have $*$: $(A \subseteq \{x. x \cdot ?v \leq 0\}) \vee A \subseteq \{x. x \cdot ?v \geq 0\}$
using *pocket-fill-line-int-aux3*[OF *assms(14) 1 2 3 4*] **by** *blast*
moreover have $q1\$2 > 0$ **using** *assms(10) assms(13) p0-def pathstart-in-path-image*
by *fastforce*
ultimately show *False*
by (*metis (no-types, lifting) IntE x assms(17) e1e2-basis(3) inner-axis*)

linorder-not-less mem-Collect-eq pathfinish-in-path-image q1-def real-inner-1-right subsetD)

qed
moreover have $q0\$2 \leq 0$ **using** *assms(10) assms(12) p1-def* **by force**
moreover have $q0 \in (\text{path-image } p) \cup (\text{path-image } q)$
by (*simp add: pathstart-in-path-image q0-def*)
ultimately show $q0\$2 = 0$ **by force**
qed

lemma *pocket-fill-line-int-aux5*:

fixes $p\ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)$ *set*
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1\$1$
defines $l \equiv \text{closed-segment } p0\ p1$
assumes *simple-path p*
assumes *simple-path q*
assumes $\text{path-image } p \cap \text{path-image } q = \{q0, q1\}$
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p\{0 < .. < 1\} \subseteq \text{interior } A$
assumes $\text{path-image } q \subseteq A$
assumes $\exists x \in p\{0 < .. < 1\}. x\$2 \geq 0$
assumes $q0 = p1 \wedge q1 = p0$
shows $l \subseteq \text{frontier } A \ \forall x \in \text{path-image } p \cup \text{path-image } q. x\$2 \geq 0$
proof–
have $1: l \subseteq \text{frontier } A$ **if** $\forall x \in \text{path-image } p \cup \text{path-image } q. x\$2 \geq 0$
proof–
have $\forall x \in \text{path-image } p \cup \text{path-image } q. x \cdot (\text{vector } [0, 1]) \geq 0$
by (*simp add: e1e2-basis(3) inner-axis that*)
then have $\forall x \in A. x \cdot (\text{vector } [0, 1]) \geq 0$
by (*smt (verit, ccfv-threshold) convex-cut-aux' assms(12) inner-commute mem-Collect-eq subset-eq*)
then have $A \subseteq \{x. x \cdot (\text{vector } [0, 1]) \geq 0\}$ **by blast**
moreover have $\text{frontier } \{x. x \cdot ((\text{vector } [0, 1]) :: (\text{real}^2)) \geq 0\} = \{x. x \cdot (\text{vector } [0, 1]) = 0\}$
by (*metis dual-order.refl frontier-halfspace-component-ge not-one-le-zero vector-2(2) zero-index*)
moreover have $l \subseteq \{x. x \cdot (\text{vector } [0, 1]) = 0\}$
proof–
have $\forall x \in l. x\$2 = 0$ **using** *assms(10) l-def segment-horizontal* **by presburger**
thus ?thesis **by** (*simp add: cart-eq-inner-axis e1e2-basis(3) subset-eq*)
qed

ultimately show *?thesis*
by (*smt (verit, best) Un-upper1 assms(12) closed-segment-subset convex-convex-hull hull-subset in-frontier-in-subset l-def p0-def p1-def pathfinish-in-path-image pathstart-in-path-image subset-eq*)
qed
have 2: *False if tht: $\neg (\forall x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 \geq 0)$*
proof –
obtain *x tx where x: $tx \in \{0..1\} \wedge q \text{ tx} = x \wedge (\forall z \in \text{path-image } p. x\$2 < z\$2)$*
using *exists-point-below-all[of p q] that*
by (*smt (verit, del-insts) tht assms(10) assms(12) assms(14) assms(7) assms(8) image-iff p0-def p1-def path-image-def pathfinish-def pathstart-def simple-path-imp-path*)
obtain *y ty where y: $ty \in \{0..1\} \wedge q \text{ ty} = y \wedge (\forall x \in \text{path-image } p. y\$2 > x\$2)$*
using *exists-point-above-all[of p q]*
by (*smt (verit, del-insts) assms(10) assms(12) assms(14) assms(16) assms(7) assms(8) image-iff p0-def p1-def path-image-def pathfinish-def pathstart-def simple-path-imp-path*)

let *?Q =*
 $\lambda q'. \text{simple-path } q' \wedge \text{path-image } p \cap \text{path-image } q' = \{ \}$
 $\wedge q' \text{ 0} = q \text{ tx} \wedge q' \text{ 1} = q \text{ ty}$
 $\wedge \text{path-image } q' \subseteq \text{path-image } q$
have *: $\bigwedge q'. ?Q \text{ } q' \implies \text{False}$
proof –
fix *q'*
assume *: *?Q q'*

have 2: *simple-path q' by (simp add: *)*
have 3: *path-image p \cap path-image q' = { } by (simp add: *)*
have 6: $\forall v \in \text{path-image } p. \text{pathstart } q' \$ 2 \leq v \$ 2$
by (*simp add: * less-eq-real-def pathstart-def x*)
have 7: $\forall v \in \text{path-image } p. v \$ 2 < \text{pathfinish } q' \$ 2$ **by** (*simp add: * pathfinish-def y*)
have 11: *path-image q' \subseteq A using * assms(15) by blast*
have $\forall x \in (\text{path-image } p) \cup (\text{path-image } q'). x\$2 \geq 0$
using *pocket-fill-line-int-aux4(2)[of p, OF - 2 3 - - 6 7 - - 11]*
by (*metis a-def assms(10) assms(11) assms(12) assms(13) assms(14) assms(7) assms(8) compact-Un compact-convex-hull compact-simple-path-image convex-convex-hull p0-def p1-def*)
thus *False*
by (*smt (verit) * UnCI assms(10) p0-def pathstart-def pathstart-in-path-image x*)
qed

have *lf: $(\forall t \in \{0..1\}. (q \text{ t} = q0 \vee q \text{ t} = q1) \longrightarrow (t = 0 \vee t = 1))$*
using *assms(8)*
unfolding *q0-def q1-def simple-path-def loop-free-def pathstart-def pathfin-*

```

ish-def
  by fastforce
  have endpoints:  $q \text{ tx} \neq q0 \wedge q \text{ ty} \neq q0 \wedge q \text{ tx} \neq q1 \wedge q \text{ ty} \neq q1$ 
  by (metis  $x \ y \ \text{assms}(10) \ \text{assms}(17) \ \text{order-less-le} \ \text{p0-def} \ \text{pathstart-in-path-image}$ )

  have  $\text{tx-neq-ty}: \text{tx} \neq \text{ty}$  using  $\text{pathstart-in-path-image} \ x \ y$  by fastforce
  moreover have  $\text{False}$  if  $\text{tx} < \text{ty}$ 
  proof-
    have  $\text{path-image } p \cap \text{path-image } (\text{subpath } \text{tx } \text{ty } q) = \{\}$ 
      (is  $\text{path-image } p \cap \text{path-image } ?q' = \{\}$ )
    proof-
      have  $q0 \notin \text{path-image } ?q' \wedge q1 \notin \text{path-image } ?q'$ 
      proof-
        have  $\{\text{tx}..\text{ty}\} \subseteq \{0..1\}$  using  $x \ y$  by simp
        then have  $(\forall t \in \{\text{tx}..\text{ty}\}. (q \ t = q0 \vee q \ t = q1) \longrightarrow (t = 0 \vee t = 1))$ 
      using  $lf$  by blast
      moreover have  $0 \notin \{\text{tx}..\text{ty}\} \wedge 1 \notin \{\text{tx}..\text{ty}\}$ 
        by (metis  $\text{atLeastAtMost-iff} \ \text{dual-order.eq-iff} \ \text{endpoints} \ \text{pathfinish-def} \ \text{pathstart-def} \ \text{q0-def} \ \text{q1-def} \ x \ y$ )
      moreover have  $\text{path-image } ?q' = q'\{\text{tx}..\text{ty}\}$  by (simp add:  $\text{path-image-subpath}$ 
        that)
      ultimately show  $?thesis$  by fastforce
    qed
  thus  $?thesis$ 
  by (smt (verit, best)  $\text{Int-empty-right} \ \text{Int-insert-right-if0} \ \text{assms}(9) \ \text{boolean-algebra-cancel.inf2} \ \text{inf.absorb-iff1} \ \text{path-image-subpath-subset} \ x \ y$ )
  qed
  thus  $?thesis$  using  $*[of \ ?q']$ 
  by (metis  $\text{assms}(8) \ \text{tx-neq-ty} \ \text{path-image-subpath-subset} \ \text{pathfinish-def} \ \text{pathfinish-subpath} \ \text{pathstart-def} \ \text{pathstart-subpath} \ \text{simple-path-subpath} \ x \ y$ )
  qed
  moreover have  $\text{False}$  if  $\text{ty} < \text{tx}$ 
  proof-
    have  $\text{path-image } p \cap \text{path-image } (\text{reversepath } (\text{subpath } \text{tx } \text{ty } q)) = \{\}$ 
      (is  $\text{path-image } p \cap \text{path-image } ?q' = \{\}$ )
    proof-
      have  $q0 \notin \text{path-image } ?q' \wedge q1 \notin \text{path-image } ?q'$ 
      proof-
        have  $\{\text{ty}..\text{tx}\} \subseteq \{0..1\}$  using  $x \ y$  by simp
        then have  $(\forall t \in \{\text{ty}..\text{tx}\}. (q \ t = q0 \vee q \ t = q1) \longrightarrow (t = 0 \vee t = 1))$ 
      using  $lf$  by blast
      moreover have  $0 \notin \{\text{ty}..\text{tx}\} \wedge 1 \notin \{\text{ty}..\text{tx}\}$ 
        by (metis  $\text{atLeastAtMost-iff} \ \text{dual-order.eq-iff} \ \text{endpoints} \ \text{pathfinish-def} \ \text{pathstart-def} \ \text{q0-def} \ \text{q1-def} \ x \ y$ )
      moreover have  $\text{path-image } ?q' = q'\{\text{ty}..\text{tx}\}$ 
        by (simp add:  $\text{path-image-subpath} \ \text{reversepath-subpath} \ \text{that}$ )
      ultimately show  $?thesis$  by fastforce
    qed
  thus  $?thesis$ 

```

```

    by (smt (verit) Int-commute assms(9) inf.absorb-iff2 inf.assoc inf-bot-right
insert-disjoint(2) path-image-reversepath path-image-subpath-subset x y)
  qed
  thus ?thesis using *[of ?q]
  by (metis * assms(8) tx-neq-ty path-image-subpath-commute path-image-subpath-subset
pathfinish-def pathfinish-subpath pathstart-def pathstart-subpath reversepath-subpath
simple-path-subpath x y)
  qed
  ultimately show False by fastforce
  qed
  show  $l \subseteq \text{frontier } A \ \forall x \in (\text{path-image } p) \cup (\text{path-image } q). \ x \geq 0$ 
  using 1 2 apply blast
  using 1 2 by blast
  qed

```

lemma *pocket-fill-line-int-aux6*:

```

fixes p q :: R-to-R2
defines p0 ≡ pathstart p
defines p1 ≡ pathfinish p
defines q0 ≡ pathstart q
defines q1 ≡ pathfinish q
defines a ≡ p1$1
assumes simple-path p
assumes simple-path q
assumes p0 = 0 ∧ p1$2 = 0
assumes a > 0
assumes q0$1 ∈ {0..a} ∧ q0$2 = 0
assumes ∀ x ∈ path-image p. q1$2 > x$2
assumes ∀ x ∈ path-image p ∪ path-image q. x$2 ≥ 0
shows path-image p ∩ path-image q ≠ {}
proof -
  let ?l1 = linepath p1 (vector [a, -1])
  let ?l2 = linepath ((vector [a, -1])::(real^2)) (vector [0, -1])
  let ?l3 = linepath ((vector [0, -1])::(real^2)) 0

  let ?R' = ?l1 +++ ?l2 +++ ?l3
  let ?R = p +++ ?R'

  have closed: closed-path ?R
  proof -
    have path ?R using assms(6) p1-def simple-path-imp-path by auto
    moreover have pathstart ?R = pathstart p by simp
    moreover have pathfinish ?R = pathfinish ?l3 by simp
    moreover have pathstart p = 0 using assms(8) p0-def by fastforce
    moreover have pathfinish ?l3 = 0 by simp
    ultimately show ?thesis unfolding closed-path-def by presburger
  qed
  have simple: simple-path ?R
  proof -

```

```

have arc ?R'
proof-
  let ?a = p1
  let ?b = (vector [a, -1])::(real^2)
  let ?c = (vector [0, -1])::(real^2)
  let ?d = 0::(real^2)

  have arcs: arc ?l1  $\wedge$  arc ?l2  $\wedge$  arc ?l3
  by (metis arc-linepath assms(8) assms(9) vector-2(1) vector-2(2) verit-comp-simplify1(1)
zero-index zero-neq-neg-one)

  have l2l3: path-image ?l2  $\cap$  path-image ?l3 = {pathfinish ?l2}
  using linepath-int-corner[of ?b ?c ?d]
  by (metis Int-commute closed-segment-commute linepath-int-corner path-image-linepath
pathfinish-linepath vector-2(2) zero-index zero-neq-neg-one)
  have l1l2: path-image ?l1  $\cap$  path-image ?l2 = {pathfinish ?l1}
  using linepath-int-corner[of ?a ?b ?c] by (simp add: assms(8))
  have l1l3: path-image ?l1  $\cap$  path-image ?l3 = {}
  using linepath-int-vertical[of ?a ?b ?c ?d] a-def assms(9) linepath-int-vertical
by auto

  have path-image ?l2  $\cap$  path-image ?l3 = {pathfinish ?l2}
  using l2l3 by blast
  moreover have sf-23: pathfinish ?l2 = pathstart ?l3 by simp
  ultimately have arc (?l2 +++ ?l3)
  by (metis arc-join-eq-alt arcs)
  moreover have path-image ?l1  $\cap$  path-image (?l2 +++ ?l3) = {pathfinish
?l1}
  using l1l2 l1l3
  by (metis (no-types, lifting) Int-Un-distrib sf-23 insert-is-Un path-image-join)
  moreover have pathfinish ?l1 = pathstart (?l2 +++ ?l3) by simp
  ultimately show arc (?l1 +++ ?l2 +++ ?l3)
  by (metis arc-join-eq-alt arcs)
qed
moreover have loop-free p using assms(6) simple-path-def by blast
moreover have path-image ?R'  $\cap$  path-image p = {p0, p1}
proof-
  have path-image ?l1  $\cap$  path-image p = {p1}
  proof-
    have  $\forall x \in \text{path-image } p. x\$2 \geq 0$  by (simp add: assms(12))
    moreover have  $\forall x \in \text{path-image } ?l1. x\$2 \leq 0$  using a-def assms(8)
segment-vertical by force
    ultimately have  $\forall x \in \text{path-image } p \cap \text{path-image } ?l1. x\$2 = 0$  by fastforce
    moreover have  $\forall x \in \text{path-image } ?l1. x\$2 = 0 \longrightarrow x = p1$ 
    by (metis (mono-tags, opaque-lifting) a-def assms(8) exhaust-2 path-image-linepath
segment-vertical vec-eq-iff vector-2(1))
    ultimately have  $\forall x \in \text{path-image } p \cap \text{path-image } ?l1. x = p1$  by fast
    moreover have  $p1 \in \text{path-image } ?l1 \wedge p1 \in \text{path-image } p$  using p1-def
by auto

```

ultimately show *?thesis by blast*
qed
moreover have *path-image ?l2* \cap *path-image p* = {}
by (*smt (verit, best) segment-horizontal assms(12) UnCI disjoint-iff path-image-linepath vector-2(2)*)
moreover have *path-image ?l3* \cap *path-image p* = {*p0*}
proof–
have $\forall x \in \text{path-image } p. x\$2 \geq 0$ **by** (*simp add: assms(12)*)
moreover have $\forall x \in \text{path-image } ?l3. x\$2 \leq 0$ **using** *a-def assms(8) segment-vertical by force*
ultimately have $\forall x \in \text{path-image } p \cap \text{path-image } ?l3. x\$2 = 0$ **by** *fastforce*
moreover have $\forall x \in \text{path-image } ?l3. x\$2 = 0 \longrightarrow x = p0$
by (*metis (no-types, opaque-lifting) assms(8) exhaust-2 path-image-linepath segment-vertical vec-eq-iff vector-2(1) zero-index*)
ultimately have $\forall x \in \text{path-image } p \cap \text{path-image } ?l3. x = p0$ **by** *fast*
moreover have *p0* \in *path-image ?l3* \wedge *p0* \in *path-image p* **using** *assms(8) p0-def by fastforce*
ultimately show *?thesis by blast*
qed
ultimately show *?thesis*
by (*smt (verit, del-insts) Int-Un-distrib Int-commute Un-assoc Un-insert-right insert-is-Un path-image-join pathfinish-linepath pathstart-join pathstart-linepath*)
qed
moreover have *arc p*
using *closed-path-def arc-distinct-ends assms(6) calculation(1) closed p1-def simple-path-imp-arc*
by *force*
ultimately show *?thesis*
by (*metis (no-types, opaque-lifting) Int-commute closed-path-def closed dual-order.refl linepath-0' p0-def p1-def pathfinish-join pathstart-def pathstart-join simple-path-join-loop-eq*)
qed

have *inside-outside: inside-outside ?R (path-inside ?R) (path-outside ?R)*
using *closed simple Jordan-inside-outside-real2*
by (*simp add: closed-path-def inside-outside-def path-inside-def path-outside-def*)

have *interior-frontier: path-inside ?R = interior (path-inside ?R)*
 \wedge *frontier (path-inside ?R) = path-image ?R*
using *inside-outside interior-open unfolding inside-outside-def by auto*

have *R-y-q1: $\forall x \in \text{path-image } ?R. x\$2 < q1\$2$*
proof–
have *: $\forall x \in \text{path-image } p. x\$2 < q1\$2$ **using** *assms(11) by blast*
moreover have $\forall x \in \text{path-image } ?l1. x\$2 < q1\$2$
using *a-def assms(8) * p1-def pathfinish-in-path-image segment-vertical by fastforce*
moreover have $\forall x \in \text{path-image } ?l2. x\$2 < q1\$2$
using *assms(8) * p1-def pathfinish-in-path-image segment-horizontal by fastforce*

moreover have $\forall x \in \text{path-image } ?l3. x\$2 < q1\$2$
using *assms(8) * p1-def pathfinish-in-path-image segment-vertical* **by** *fastforce*
ultimately show *?thesis* **by** (*metis not-in-path-image-join*)
qed
have *R-y-0: $\forall x \in \text{path-image } ?R. x\$2 \geq -1$*
proof–
have $\forall x \in \text{path-image } ?l1. x\$2 \geq -1$ **using** *a-def assms(8) segment-vertical*
by *fastforce*
moreover have $\forall x \in \text{path-image } ?l2. x\$2 \geq -1$ **using** *segment-horizontal* **by**
auto
moreover have $\forall x \in \text{path-image } ?l3. x\$2 \geq -1$ **using** *segment-vertical* **by**
auto
moreover have $\forall x \in \text{path-image } p. x\$2 \geq -1$ **using** *assms(12)* **by** *force*
ultimately show *?thesis* **by** (*metis not-in-path-image-join*)
qed

have *?thesis* **if** $p0 \in \text{path-image } q \vee p1 \in \text{path-image } q$ **using** *p0-def p1-def* **that**
by *blast*
moreover have *?thesis* **if** $p0 \notin \text{path-image } q \wedge p1 \notin \text{path-image } q \wedge q0 \notin$
path-image } p
proof–
have *q-int-l1: $\text{path-image } q \cap \text{path-image } ?l1 = \{\}$*
proof–
have $\forall x \in \text{path-image } q. x\$2 \geq 0$ **by** (*simp add: assms(12)*)
moreover have $\forall x \in \text{path-image } ?l1. x\$2 = 0 \longrightarrow x = p1$
by (*metis (mono-tags, opaque-lifting) a-def assms(8) exhaust-2 path-image-linepath*
segment-vertical vec-eq-iff vector-2(1))
ultimately show *?thesis* **using** *that a-def assms(8) segment-vertical* **by**
fastforce
qed
moreover have *q-int-l2: $\text{path-image } q \cap \text{path-image } ?l2 = \{\}$*
by (*smt (verit, ccfv-threshold) UnCI assms(12) disjoint-iff path-image-linepath*
segment-horizontal vector-2(2))
moreover have *q-int-l3: $\text{path-image } q \cap \text{path-image } ?l3 = \{\}$*
proof–
have $\forall x \in \text{path-image } q. x\$2 \geq 0$ **by** (*simp add: assms(12)*)
moreover have $\forall x \in \text{path-image } ?l3. x\$2 = 0 \longrightarrow x = p0$
by (*metis (no-types, opaque-lifting) assms(8) exhaust-2 path-image-linepath*
segment-vertical vec-eq-iff vector-2(1) zero-index)
ultimately show *?thesis* **using** *that a-def assms(8) segment-vertical* **by**
fastforce
qed
ultimately have *q0-notin-R: $q0 \notin \text{path-image } ?R$*
using *that* **by** (*simp add: disjoint-iff not-in-path-image-join pathstart-in-path-image*
q0-def)

have $\text{path-image } q \cap \text{path-image } ?R \neq \{\}$
proof–
have $q0 \in \text{path-inside } ?R$

proof-
let $?e = (\text{vector } [q0\$1, -1])::(\text{real}^2)$
let $?d1 = (\text{vector } [a, -1])::(\text{real}^2)$
let $?d2 = (\text{vector } [0, -1])::(\text{real}^2)$

have $0 < q0\$1 \wedge q0\$1 < a$
by (*smt (verit) a-def assms(10) assms(8) atLeastAtMost-iff exhaust-2 linorder-not-less pathstart-in-path-image q0-def that vec-eq-iff zero-index*)
then have $q0\$1 > 0 \wedge a - q0\$1 > 0$ **by** *simp*
then have $\min(\min(q0\$1) (a - q0\$1)) > 0$ (**is** $?e' > 0$) **by** *linarith*
then have $0 < ?e'/2 \wedge ?e'/2 < 1 \wedge ?e'/2 < q0\$1 \wedge ?e'/2 < a - q0\$1$

by *argo*
then obtain ε **where** $\varepsilon: 0 < \varepsilon \wedge \varepsilon < 1 \wedge \varepsilon < q0\$1 \wedge \varepsilon < a - q0\1 **by**
blast

moreover have $?e \in \text{frontier } (\text{path-inside } ?R)$
by (*smt (verit, del-insts) UnCI <0 < q0 \$ 1 \wedge 0 < a - q0 \$ 1> interior-frontier p1-def path-image-join path-image-linepath pathfinish-linepath pathstart-join pathstart-linepath segment-horizontal vector-2(1) vector-2(2)*)
ultimately obtain *int-p* **where** *int-p*: $\text{int-p} \in \text{ball } ?e \varepsilon \cap \text{path-inside } ?R$
by (*meson inside-outside frontier-straddle mem-ball IntI*)

have *int-p*: $\text{int-p}\$1 > 0 \wedge \text{int-p}\$1 < a$

proof-
have $\text{int-p}\$1 > 0$
proof(*rule ccontr*)
assume $\neg \text{int-p}\$1 > 0$
moreover have $\text{dist } (\text{int-p}\$1) (q0\$1) < q0\1
by (*smt (verit) IntE \varepsilon dist-commute dist-vec-nth-le int-p mem-ball vector-2(1)*)
ultimately show *False* **using** *dist-real-def* **by** *force*
qed

moreover have $\text{int-p}\$1 < a$
proof(*rule ccontr*)
assume $\neg \text{int-p}\$1 < a$
moreover have $\text{dist } (\text{int-p}\$1) (q0\$1) < a - q0\1
by (*smt (verit) IntE \varepsilon dist-commute dist-vec-nth-le int-p mem-ball vector-2(1)*)
ultimately show *False* **using** *dist-real-def* **by** *force*
qed

ultimately show *?thesis* **by** *blast*

qed

have *int-p*: $\text{int-p}\$2 > -1 \wedge \text{int-p}\$2 < 0$

proof-
have $\text{int-p}\$2 > -1$
proof(*rule ccontr*)
assume $\neg \text{int-p}\$2 > -1$
then have $\text{int-p}\$2 \leq -1$ **by** *simp*
let $?e2' = (\text{vector } [0, -1])::(\text{real}^2)$
let $?ray = \lambda d. \text{int-p} + d *_R ?e2'$

```

    have  $\neg (\exists d > 0. \text{?ray } d \in \text{path-image } ?R)$ 
  proof-
    have  $\forall d > 0. (\text{?ray } d)\$2 < -1$  using * by auto
    thus ?thesis using R-y-0 by force
  qed
  moreover have bounded (path-inside ?R) using bounded-finite-inside
simple by blast
  moreover have  $\text{?e2}' \neq 0$  by (metis vector-2(2) zero-index zero-neq-neg-one)
  ultimately have  $\text{int-p} \notin \text{path-inside } ?R$ 
    using ray-to-frontier[of path-inside ?R] interior-frontier by metis
  thus False using int-p by blast
  qed
  moreover have  $\text{int-p}\$2 < 0$ 
  proof(rule ccontr)
    assume  $\neg \text{int-p}\$2 < 0$ 
    then have  $\text{dist int-p } ?e \geq 1$ 
      by (smt (verit, del-insts) dist-real-def dist-vec-nth-le vector-2(2))
  thus False by (smt (verit, del-insts) IntD1  $\varepsilon$  dist-commute int-p mem-ball)
  qed
  ultimately show ?thesis by blast
  qed

let ?int-l = linepath int-p q0

have path-image ?int-l  $\cap$  path-image ?l1 = {}
  using  $\langle 0 < q0 \$ 1 \wedge q0 \$ 1 < a \rangle$  a-def int-p-x linepath-int-columns by
auto
  moreover have path-image ?int-l  $\cap$  path-image ?l2 = {}
    by (smt (verit, best) assms(10) disjoint-iff int-p-y linepath-int-rows vec-
tor-2(2))
  moreover have path-image ?int-l  $\cap$  path-image ?l3 = {}
    by (smt (verit, del-insts)  $\varepsilon$  disjoint-iff int-p-x linepath-int-columns vec-
tor-2(1) zero-index)
  moreover have path-image ?int-l  $\cap$  path-image p = {}
  proof-
    have  $\forall t \in \{0..1\}. (\text{?int-l } t)\$2 = 0 \longrightarrow t = 1$ 
      unfolding linepath-def using assms(10) int-p-y by force
    then have  $\forall x \in \text{path-image } ?int-l. x\$2 = 0 \longrightarrow x = q0$ 
      unfolding path-image-def using linepath-1' by fastforce
    moreover have  $\forall x \in \text{path-image } p. x\$2 \geq 0$  by (simp add: assms(12))
    moreover have  $\forall x \in \text{path-image } ?int-l. x\$2 \leq 0$ 
      by (smt (verit) assms(10) int-p-y linepath-bound-2(2))
    ultimately show ?thesis using that by fastforce
  qed
  ultimately have path-image ?int-l  $\cap$  path-image ?R = {}
    by (simp add: disjoint-iff not-in-path-image-join)

    then have path-image ?int-l  $\subseteq$  path-inside ?R  $\vee$  path-image ?int-l  $\subseteq$ 
path-outside ?R

```

by (*metis IntD2 IntI convex-imp-path-connected convex-segment(1) empty-iff
int-p interior-frontier path-connected-not-frontier-subset path-image-linepath path-
start-in-path-image pathstart-linepath*)

moreover have $?int-l\ 0 = int-p \wedge int-p \in path-inside\ ?R$
using *int-p* **by** (*simp add: linepath-0'*)
ultimately have $path-image\ ?int-l \subseteq path-inside\ ?R$
using *inside-outside-def local.inside-outside* **by auto**
thus *?thesis* **by auto**

qed
then have $q0 \in -\ (path-outside\ ?R)$
by (*metis ComplI IntI equals0D inside-Int-outside path-inside-def path-outside-def*)
moreover have $q1 \in path-outside\ ?R$
proof-
let $?e2 = (vector\ [0,\ 1])::(real^2)$
let $?ray = \lambda d. q1 + d *R\ ?e2$
have $\neg (\exists d > 0. ?ray\ d \in path-image\ ?R)$
proof-
have $\forall d > 0. (?ray\ d)\$2 > q1\$2$ **by simp**
thus *?thesis* **using** *R-y-q1* **by fastforce**
qed
moreover have *bounded (path-inside ?R) using bounded-finite-inside simple*

by blast
moreover have $?e2 \neq 0$ **using** *e1e2-basis(4)* **by force**
ultimately have $q1 \notin path-inside\ ?R$
using *ray-to-frontier[of path-inside ?R] interior-frontier* **by metis**
moreover have $q1 \notin path-image\ ?R$ **using** *R-y-q1* **by blast**
ultimately show *?thesis* **using** *inside-outside unfolding inside-outside-def*

by blast
qed
ultimately have $path-image\ q \cap -\ (path-outside\ ?R) \neq \{\}$
 $\wedge path-image\ q \cap (path-outside\ ?R) \neq \{\}$
using *q0-def q1-def* **by blast**
moreover have *path-connected (path-image q)*
using *assms(7) path-connected-path-image simple-path-def* **by blast**
moreover have $path-image\ ?R = frontier\ (path-outside\ ?R)$
using *inside-outside unfolding inside-outside-def p0-def path-inside-def* **by**

blast
ultimately show *?thesis* **by** (*metis Diff-eq Diff-eq-empty-iff path-connected-not-frontier-subset*)
qed
thus *?thesis* **by** (*meson q-int-l1 q-int-l2 q-int-l3 disjoint-iff not-in-path-image-join*)
qed
ultimately show *?thesis* **using** *q0-def* **by blast**

qed

lemma *pocket-fill-line-int-aux7:*
fixes $p\ q :: R\ to\ R2$
fixes $A :: (real^2)\ set$
defines $p0 \equiv pathstart\ p$
defines $p1 \equiv pathfinish\ p$

```

defines  $q0 \equiv \text{pathstart } q$ 
defines  $q1 \equiv \text{pathfinish } q$ 
defines  $a \equiv p1\$1$ 
defines  $l \equiv \text{open-segment } p0 \ p1$ 
assumes simple-path  $p$ 
assumes simple-path  $q$ 
assumes  $\text{path-image } p \cap \text{path-image } q = \{q0, q1\}$ 
assumes  $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$ 
assumes  $a > 0$ 
assumes  $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$ 
assumes  $\{p0, p1\} \subseteq \text{frontier } A$ 
assumes  $p\{0<..<<1\} \subseteq \text{interior } A$ 
assumes  $\exists x \in p\{0<..<<1\}. x\$2 \geq 0$ 
assumes  $q0 = p1 \wedge q1 = p0$ 
shows  $\text{path-image } q \cap l = \{\}$  closed-segment  $p0 \ p1 \subseteq \text{frontier } A$ 
proof –
  have 1:  $\text{path-image } p \cap \text{path-image } q = \{\text{pathstart } q, \text{pathfinish } q\}$ 
    by (simp add: assms(9) q0-def q1-def)
  have 2:  $\text{pathstart } p \ \$1 = 0 \wedge \text{pathstart } p \ \$2 = 0 \wedge \text{pathfinish } p \ \$2 = 0$ 
    using assms(10) p0-def p1-def by blast
  have 3:  $0 < \text{pathfinish } p \ \$1$  using a-def assms(11) p1-def by auto
  have 4:  $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$  by (simp add: assms(12))
  have 5:  $\{\text{pathstart } p, \text{pathfinish } p\} \subseteq \text{frontier } A$  using assms(13) p0-def p1-def
by blast
  have 6:  $p \ \{0<..<<1\} \subseteq \text{interior } A$  using assms(14) by blast
  have 7:  $\text{path-image } q \subseteq A$  using assms(12) hull-subset by force
  have 8:  $\exists x \in p\{0<..<<1\}. x\$2 \geq 0$  using assms(15) by blast
  have 9:  $\text{pathstart } q = \text{pathfinish } p \wedge \text{pathfinish } q = \text{pathstart } p$ 
    using assms(16) p0-def p1-def q0-def q1-def by fastforce
  have *:  $\forall x \in (\text{path-image } p) \cup (\text{path-image } q). x\$2 \geq 0$ 
    using pocket-fill-line-int-aux5(2)[OF assms(7) assms(8) 1 2 3 4 5 6 7 8 9] by
blast

show  $\text{closed-segment } p0 \ p1 \subseteq \text{frontier } A$ 
  using pocket-fill-line-int-aux5(1)[OF assms(7) assms(8) 1 2 3 4 5 6 7 8 9]
  unfolding l-def p0-def p1-def by blast
show  $\text{path-image } q \cap l = \{\}$ 
proof(rule ccontr)
  assume  $\neg \text{path-image } q \cap l = \{\}$ 
  then obtain  $x \ tx$  where  $x: tx \in \{0..1\} \wedge q \ tx = x \wedge x \in l$ 
    by (metis (no-types, lifting) disjoint-iff imageE path-image-def)
  obtain  $y \ ty$  where  $y: ty \in \{0..1\} \wedge q \ ty = y \wedge (\forall x \in \text{path-image } p. y\$2 >$ 
 $x\$2)$ 
    using exists-point-above-all[of p q]
    by (smt (verit, del-insts) 4 6 8 assms(10) assms(7) assms(8) p0-def p1-def
pathfinish-def pathstart-def simple-path-def image-iff path-image-def)

  have lf:  $(\forall t \in \{0..1\}. (q \ t = q0 \vee q \ t = q1) \longrightarrow (t = 0 \vee t = 1))$ 
    using assms(8)

```

```

    unfolding q0-def q1-def simple-path-def loop-free-def pathstart-def pathfin-
ish-def
    by fastforce
    have endpoints: q tx ≠ q0 ∧ q ty ≠ q0 ∧ q tx ≠ q1 ∧ q ty ≠ q1 ∧ tx ≠ ty
    proof-
    have (q ty)$2 > 0 by (metis assms(10) p0-def pathstart-in-path-image y)
    moreover have (q tx)$2 = 0
    proof-
    have q tx ∈ closed-segment q0 q1
    using assms(16) l-def open-closed-segment open-segment-commute x by
blast
    thus ?thesis by (simp add: assms(10) assms(16) segment-horizontal)
    qed
    moreover have q0 ∉ open-segment q0 q1 ∧ q1 ∉ open-segment q0 q1
    by (simp add: open-segment-def)
    ultimately show ?thesis
    using assms(10) assms(16) l-def open-segment-commute x by auto
    qed

let ?Q =
  λq'. simple-path q' ∧ path-image p ∩ path-image q' = {}
  ∧ q' 0 = q tx ∧ q' 1 = q ty
  ∧ path-image q' ⊆ path-image q
have **: ∧q'. ?Q q' ⇒ False
proof-
fix q'
assume **: ?Q q'
have 1: simple-path q' by (simp add: **)
have 2: pathstart p = 0 ∧ pathfinish p $ 2 = 0
by (metis (mono-tags, lifting) assms(10) exhaust-2 p0-def p1-def vec-eq-iff
zero-index)
have 3: 0 < pathfinish p $ 1 using a-def assms(11) p1-def by blast
have 4: pathstart q' $ 1 ∈ {0..pathfinish p $ 1} ∧ pathstart q' $ 2 = 0
proof-
have q' 0 ∈ closed-segment p0 p1 using ** l-def open-closed-segment x by
auto
thus ?thesis
by (smt (z3) 2 a-def assms(11) atLeastAtMost-iff atLeastatMost-empty
p0-def p1-def pathstart-def pathstart-subpath segment-horizontal zero-index)
qed
have 5: ∀x∈path-image p. x $ 2 < pathfinish q' $ 2 by (simp add: **
pathfinish-def y)
have 6: ∀x∈path-image p ∪ path-image q'. 0 ≤ x $ 2 using * ** by blast
have path-image p ∩ path-image q' ≠ {}
using pocket-fill-line-int-aux6[OF assms(7) 1 2 3 4 5 6] by simp
thus False using ** by blast
qed

have False if tx < ty

```

```

proof-
  let ?q' = subpath tx ty q
  have q0 ∉ path-image ?q' ∧ q1 ∉ path-image ?q'
  proof-
    have {tx..ty} ⊆ {0..1} using x y by simp
    then have (∀ t ∈ {tx..ty}. (q t = q0 ∨ q t = q1) → (t = 0 ∨ t = 1))
using lf by blast
    moreover have 0 ∉ {tx..ty} ∧ 1 ∉ {tx..ty}
      by (metis atLeastAtMost-iff dual-order.eq-iff endpoints pathfinish-def
pathstart-def q0-def q1-def x y)
    moreover have path-image ?q' = q'{tx..ty} by (simp add: path-image-subpath
that)
    ultimately show ?thesis by fastforce
  qed
  then have ?Q ?q'
    by (smt (verit, best) assms(8) assms(9) disjoint-insert(1) endpoints
inf.absorb-iff1 inf-bot-right inf-left-commute path-image-subpath-subset pathfinish-def
pathfinish-subpath pathstart-def pathstart-subpath simple-path-subpath x y)
  thus False using ** by auto
  qed
  moreover have False if tx > ty
  proof-
    let ?q' = reversepath (subpath ty tx q)
    have q0 ∉ path-image ?q' ∧ q1 ∉ path-image ?q'
    proof-
      have {ty..tx} ⊆ {0..1} using x y by simp
      then have (∀ t ∈ {ty..tx}. (q t = q0 ∨ q t = q1) → (t = 0 ∨ t = 1))
using lf by blast
      moreover have 0 ∉ {ty..tx} ∧ 1 ∉ {ty..tx}
        by (metis atLeastAtMost-iff dual-order.eq-iff endpoints pathfinish-def
pathstart-def q0-def q1-def x y)
      moreover have path-image ?q' = q'{ty..tx} by (simp add: path-image-subpath
that)
      ultimately show ?thesis by fastforce
    qed
    then have ?Q ?q'
      by (smt (verit) assms(8) assms(9) endpoints inf.absorb-iff2 inf.assoc
inf-bot-left insert-disjoint(2) path-image-subpath-subset pathstart-def pathstart-subpath
reversepath-def reversepath-subpath simple-path-subpath x y)
    thus False using ** by blast
  qed
  ultimately show False using endpoints by linarith
  qed
  qed

```

lemma frontier-injective-linear-image:
fixes f :: 'a::euclidean-space ⇒ 'a::euclidean-space
assumes linear f inj f

shows $f'(\text{frontier } S) = \text{frontier } (f' S)$
using *interior-injective-linear-image closure-injective-linear-image frontier-def*
assms
by (*metis image-set-diff*)

lemma *pocket-fill-line-int-aux8*:

fixes $p\ q :: R\text{-to-}R^2$
fixes $A :: (\text{real}^2)\ \text{set}$
defines $p0 \equiv \text{pathstart } p$
defines $p1 \equiv \text{pathfinish } p$
defines $q0 \equiv \text{pathstart } q$
defines $q1 \equiv \text{pathfinish } q$
defines $a \equiv p1\$1$
defines $l \equiv \text{open-segment } p0\ p1$
assumes *simple-path* p
assumes *simple-path* q
assumes $\text{path-image } p \cap \text{path-image } q = \{q0, q1\}$
assumes $p0\$1 = 0 \wedge p0\$2 = 0 \wedge p1\$2 = 0$
assumes $a > 0$
assumes $A = \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
assumes $\{p0, p1\} \subseteq \text{frontier } A$
assumes $p'\{0 < .. < 1\} \subseteq \text{interior } A$
assumes $q0 = p1 \wedge q1 = p0$
shows $\text{path-image } q \cap l = \{\} \wedge l \subseteq \text{frontier } A$
proof –
have *?thesis* **if** $ex: \exists x \in p'\{0 < .. < 1\}. x\$2 \geq 0$
using $ex\ a\text{-def}\ \text{assms}\ \text{dual-order.trans}\ l\text{-def}\ p0\text{-def}\ p1\text{-def}\ \text{pocket-fill-line-int-aux7}(1)$
pocket-fill-line-int-aux7(2) q0-def q1-def segment-open-subset-closed **that**

by (*smt* (*verit*) $a\text{-def}\ \text{assms}\ \text{dual-order.trans}\ l\text{-def}\ p0\text{-def}\ p1\text{-def}\ \text{pocket-fill-line-int-aux7}(1)$
pocket-fill-line-int-aux7(2) q0-def q1-def segment-open-subset-closed **that**)
moreover **have** *?thesis* **if** $\neg (\exists x \in p'\{0 < .. < 1\}. x\$2 \geq 0)$

proof –
let $?M = (\text{vector } [\text{vector } [1, 0], \text{vector } [0, -1]]) :: (\text{real}^2 \Rightarrow \text{real}^2)$
let $?f = \lambda v. ?M * v$
let $?g = (\lambda v. \text{vector } [v\$1, -v\$2]) :: (\text{real}^2 \Rightarrow \text{real}^2)$
define p' **where** $p' \equiv ?f \circ p$
define q' **where** $q' \equiv ?f \circ q$
define A' **where** $A' \equiv ?f' A$

have *inj*: *inj* $?f$ **and** *f-eq-g*: $?f = ?g$
using *flip-function*(1) **apply** *blast*
using *flip-function*(2) **by** *blast*

have 4 : $\text{pathstart } p'\$1 = 0 \wedge \text{pathstart } p'\$2 = 0 \wedge \text{pathfinish } p'\$2 = 0$
by (*smt* (*verit*, *best*) $\text{assms}(10)\ f\text{-eq-g}\ o\text{-apply}\ p'\text{-def}\ p0\text{-def}\ p1\text{-def}\ \text{pathfinish-def}\ \text{pathstart-def}\ \text{vector-2}(1)\ \text{vector-2}(2)$)
have *startfinish*: $\text{pathstart } p' = \text{pathstart } p \wedge \text{pathfinish } p' = \text{pathfinish } p$
by (*metis* (*mono-tags*, *opaque-lifting*) $4\ \text{assms}(10)\ \text{exhaust-2}\ f\text{-eq-g}\ o\text{-apply}$)

p' -def $p0$ -def $p1$ -def $pathfinish$ -def vec -eq-iff $vector$ -2(1))

have 1: *simple-path* p' **using** *inj* **by** (*simp add: assms(7) simple-path-linear-image-eq* p' -def)

have 2: *simple-path* q' **using** *inj* **by** (*simp add: assms(8) simple-path-linear-image-eq* q' -def)

have 3: $path\text{-}image\ p' \cap path\text{-}image\ q' = \{pathstart\ q', pathfinish\ q'\}$

proof–

have $path\text{-}image\ p' \cap path\text{-}image\ q' = ?f'(path\text{-}image\ p \cap path\text{-}image\ q)$

unfolding p' -def q' -def **by** (*simp add: image-Int inj path-image-compose*)

also have $\dots = ?f'\{q0, q1\}$ **using** *assms(9)* **by** *presburger*

finally show *?thesis*

by (*simp add: startfinish pathfinish-compose pathstart-compose* q' -def $q0$ -def $q1$ -def)

qed

have 5: $0 < pathfinish\ p' \ \$\ 1$

by (*metis (mono-tags, lifting) a-def assms(11) f-eq-g o-apply* p' -def $p1$ -def $pathfinish$ -def $vector$ -2(1))

have 6: $A' = convex\ hull\ (path\text{-}image\ p' \cup path\text{-}image\ q')$

proof–

have $path\text{-}image\ (?f \circ p) = ?f'(path\text{-}image\ p)$ **using** *path-image-compose* **by** *blast*

moreover have $path\text{-}image\ (?f \circ q) = ?f'(path\text{-}image\ q)$ **using** *path-image-compose* **by** *blast*

moreover have $?f'(path\text{-}image\ p \cup path\text{-}image\ q) = ?f'(path\text{-}image\ p) \cup ?f'(path\text{-}image\ q)$

by *blast*

moreover have $A' = convex\ hull\ (?f'(path\text{-}image\ p \cup path\text{-}image\ q))$

by (*simp add: assms(12) convex-hull-linear-image* A' -def)

ultimately show *?thesis* **using** p' -def q' -def A' -def **by** *argo*

qed

have 7: $\{pathstart\ p', pathfinish\ p'\} \subseteq frontier\ A'$

using *frontier-injective-linear-image*

by (*smt (verit, best) $\exists A'$ -def assms(13) assms(15) assms(9) doubleton-eq-iff image-Int inj inj-image-subset-iff matrix-vector-mul-linear* p' -def $p0$ -def $p1$ -def $path\text{-}image\text{-}linear\text{-}image\ pathfinish\text{-}compose\ pathstart\text{-}compose\ q'$ -def $q0$ -def $q1$ -def)

have 8: $p'\{0 < .. < 1\} \subseteq interior\ A'$

proof–

have $?f'(interior\ A) = interior\ A'$ **by** (*simp add: A'-def inj interior-injective-linear-image*)

thus *?thesis* **using** *assms(14)* p' -def **by** *auto*

qed

have 9: $\exists x \in p'\{0 < .. < 1\}. x\$2 \geq 0$

proof–

have $\exists x \in p'\{0 < .. < 1\}. x\$2 < 0$

by (*metis that all-not-in-conv bot.extremum greaterThanLessThan-subseteq-greaterThanLessThan image-is-empty verit-comp-simplify1* (3) *zero-less-one*)

then obtain x **where** $x \in p'\{0 < .. < 1\} \wedge x\$2 < 0$ **by** *presburger*

moreover then have $(?g\ x)\$2 > 0$ **by** *fastforce*

ultimately show *?thesis* **by** (*smt (verit, ccfv-threshold) f-eq-g image-iff*)

o-apply p'-def)
qed
have 10: *pathstart q' = pathfinish p' \wedge pathfinish q' = pathstart p'*
by (*metis (mono-tags, lifting) assms(15) o-apply p'-def p0-def p1-def pathfin-*
ish-def pathstart-def q'-def q0-def q1-def)

have *path-image q' \cap open-segment (pathstart p') (pathfinish p') = {}*
using *pocket-fill-line-int-ax7(1)[OF 1 2 3 4 5 6 7 8 9 10]* **by** *blast*
then have *path-image q' \cap l = {}* **using** *startfinish unfolding l-def p0-def*
p1-def **by** *simp*
moreover have *on-l: $\bigwedge x. x \in l \implies ?g x \in l$*
proof–
fix *x :: real²*
assume *x \in l*
moreover then have *x² = 0* **by** (*metis assms(6,10) segment-horizontal*
open-closed-segment)
moreover then have *(?g x)² = 0* **by** *simp*
moreover have *(?g x)¹ = x¹* **by** *simp*
ultimately show *?g x \in l* **by** (*smt (verit, ccfv-SIG) exhaust-2 vec-eq-iff*)
qed
ultimately have *path-image q \cap l = {}*
by (*metis (no-types, lifting) disjoint-iff f-eq-g image-eqI path-image-compose*
q'-def)
moreover have *l \subseteq frontier A*
proof–
have *pathstart p' = pathstart p \wedge pathfinish p' = pathfinish p*
using *startfinish* **by** *auto*
then have *?f'l \subseteq frontier A'*
using *pocket-fill-line-int-ax7(2)[OF 1 2 3 4 5 6 7 8 9 10]* *on-l f-eq-g l-def*
p0-def p1-def segment-open-subset-closed
by *force*
thus *?thesis*
by (*metis (no-types, lifting) A'-def frontier-injective-linear-image inj inj-image-subset-iff*
matrix-vector-mul-linear)
qed
ultimately show *?thesis* **by** *fast*
qed
ultimately show *?thesis* **by** *argo*
qed

lemma *simple-path-linear-image:*

assumes *simple-path p*
assumes *inj f \wedge bounded-linear f*
shows *simple-path (f \circ p)*

proof–

have *continuous-on {x. True} f* **using** *assms(2) linear-continuous-on* **by** *blast*

then have 1: *path (f \circ p)*

by (*metis Collect-cong UNIV-I assms(1) continuous-on-subset path-continuous-image*
simple-path-imp-path top-empty-eq top-greatest top-set-def)

have *inj-on* $p \{0 < .. < 1\}$ **by** (*simp add: assms(1) simple-path-inj-on*)
then have *inj-on* $(f \circ p) \{0 < .. < 1\}$ **by** (*meson assms(2) comp-inj-on inj-on-subset top-greatest*)
then have *loop-free* $(f \circ p)$
by (*metis (mono-tags, lifting) assms(1) assms(2) comp-apply inj-eq loop-free-def simple-path-def*)
thus *?thesis* **using** 1 **unfolding** *simple-path-def* **by** *blast*
qed

lemma *pts-interior*:

fixes *pts*
defines $p \equiv \text{make-polygonal-path } pts$
assumes *convex* H
assumes $\forall j \in \{0 < .. < \text{length } pts - 1\}. pts!j \notin \text{frontier } H$
assumes *loop-free* p
assumes *path-image* $p \subseteq H$
assumes $\text{length } pts \geq 3$
shows $p\{0 < .. < 1\} \subseteq \text{interior } H$
proof(*rule subsetI*)
fix x **assume** $*$: $x \in p\{0 < .. < 1\}$
then obtain t **where** $t: x = p \ t \wedge t \in \{0 < .. < 1\}$ **by** *blast*
then have $x \neq p \ 0 \wedge x \neq p \ 1$ **using** *assms(4)* **unfolding** *loop-free-def* **by** *fastforce*
then have *x-neq*: $x \neq \text{hd } pts \wedge x \neq \text{last } pts$
by (*metis assms(4) constant-linepath-is-not-loop-free hd-conv-nth last-conv-nth make-polygonal-path.simps(1) p-def pathfinish-def pathstart-def polygon-pathfinish polygon-pathstart*)

have $x \in \text{interior } H$ **if** $*$: $\exists i < \text{length } pts. x = pts!i$

proof–

obtain i **where** $i: i < \text{length } pts \wedge x = pts!i$ **using** $*$ **by** *blast*

then have $i \neq 0 \wedge i \neq \text{length } pts - 1$

by (*metis x-neq gr-implies-not0 hd-conv-nth last-conv-nth list.size(3)*)

then have $i \in \{0 < .. < \text{length } pts - 1\}$ **using** i **by** *fastforce*

then have $pts!i \notin \text{frontier } H$ **using** *assms(3)* **by** *blast*

then have $pts!i \in \text{interior } H$

by (*metis DiffI assms(5) closure-subset frontier-def i nth-mem p-def subsetD vertices-on-path-image*)

thus *?thesis* **using** *assms(3)* i **by** *blast*

qed

moreover have $x \in \text{interior } H$ **if** $*$: $\neg (\exists i < \text{length } pts. x = pts!i)$

proof–

have $x \in \text{path-image } p$ **using** $*$ **unfolding** *path-image-def* **by** *force*

then obtain i **where** $i: x \in \text{path-image } (\text{linepath } (pts!i) (pts!(i+1))) \wedge i < \text{length } pts - 1$

using *make-polygonal-path-image-property[of pts x] assms(6)* **unfolding** *p-def* **by** *auto*

moreover then have $x \neq pts!i \wedge x \neq pts!(i+1)$ **using** $*$ **by** *force*

ultimately have $x \in \text{open-segment } (vts!i) (vts!(i+1))$ **by** (*simp add: open-segment-def*)
moreover then have $x \in \text{rel-interior } (\text{path-image } (\text{linepath } (vts!i) (vts!(i+1))))$
by (*metis empty-iff open-segment-idem path-image-linepath rel-interior-closed-segment*)
moreover have *interior-nonempty: vts!i ∈ interior H ∨ vts!(i+1) ∈ interior*
H
proof(*rule ccontr*)
assume $\neg (vts!i \in \text{interior } H \vee vts!(i+1) \in \text{interior } H)$
then have $vts!i \in \text{frontier } H \wedge vts!(i+1) \in \text{frontier } H$
using *assms(5) closure-subset frontier-def i p-def vertices-on-path-image* **by**
fastforce
thus *False*
by (*metis assms(3) i Suc-1 Suc-eq-plus1 add commute add.right-neutral*
assms(6) eval-nat-numeral(3) greaterThanLessThan-iff less-diff-conv linorder-not-le
not-gr-zero not-less-eq-eq)
qed
ultimately have $x \in \text{rel-interior } H$
by (*smt (verit, ccfv-SIG) add-diff-inverse-nat assms(2) assms(5) convex-same-rel-interior-closure-straddle*
empty-iff i in-interior-closure-convex-segment less-diff-conv less-nat-zero-code nat-diff-split
nth-mem open-segment-commute p-def rel-interior-nonempty-interior subset-eq trans-less-add2
vertices-on-path-image)
moreover have $\text{interior } H \neq \{\}$ **using** *interior-nonempty* **by** *blast*
ultimately show *?thesis* **using** *rel-interior-nonempty-interior* **by** *blast*
qed
ultimately show $x \in \text{interior } H$ **by** *blast*
qed

lemma *pocket-fill-line-int-0:*
assumes *polygon-of r vts*
defines $H \equiv \text{convex hull } (\text{set } vts)$
assumes $2 \leq i \wedge i < \text{length } vts - 1$
defines $a \equiv \text{hd } vts$
defines $b \equiv vts!i$
assumes $\{a, b\} \subseteq \text{frontier } H$
assumes $\forall j \in \{0 <.. <i\}. vts!j \notin \text{frontier } H$
assumes $a = 0$
shows $\text{path-image } (\text{linepath } a b) \cap \text{path-image } r = \{a, b\}$
 $\text{path-image } (\text{linepath } a b) \subseteq \text{frontier } H$
proof–
let $?x = (b - a)$
let $?e = \text{norm } (b - a) *_R ((\text{vector } [1, 0])::(\text{real}^2))$
have $\text{norm } ?x = \text{norm } ?e$ **by** (*simp add: e1e2-basis(1)*)
then obtain f **where** $f: \text{orthogonal-transformation } f \wedge \det(\text{matrix } f) = 1 \wedge f$
 $?x = ?e$
using *rotation-exists* **by** (*metis two-le-card*)

have $\text{bij: } \text{bij } f \wedge \text{linear } f$
using *f orthogonal-transformation-bij orthogonal-transformation-def* **by** *blast*

let $?p\text{-}vts = \text{take } (i + 1) vts$

```

let ?q-vts = drop i vts
let ?p = make-polygonal-path ?p-vts
let ?q = make-polygonal-path ?q-vts

let ?p' = f ∘ ?p
let ?q' = f ∘ ?q
let ?H' = convex hull (path-image ?p' ∪ path-image ?q')

have vts-split: vts = ?p-vts @ (tl ?q-vts)
  by (metis Suc-eq-plus1 append-take-drop-id drop-Suc tl-drop)

have simple-path r using assms(1) unfolding polygon-of-def polygon-def by
blast
  then have a-neq-b: a ≠ b
    using simple-polygonal-path-vts-distinct[of vts]
    by (metis (mono-tags, lifting) a-def assms(1) assms(3) b-def bot-nat-0.extremum-strict
butlast-conv-take constant-linepath-is-not-loop-free distinct-nth-eq-iff dual-order.strict-trans2
hd-conv-nth length-butlast make-polygonal-path.simps(1) nat-neq-iff nth-take poly-
gon-of-def pos2 simple-path-def)

  have H-r: H = convex hull (path-image r)
    by (metis (no-types, lifting) H-def Un-subset-iff assms(1) convex-convex-hull
convex-hull-eq convex-hull-of-polygon-is-convex-hull-of-vts hull-mono hull-subset or-
der-antisym-conv polygon-of-def vertices-on-path-image)
    moreover have r-union: path-image r = (path-image ?p) ∪ (path-image ?q)
    proof –
      let ?i = i + 1
      let ?x = ((2::real) ^ (?i - 1) - 1) / 2 ^ (?i - 1)
      have ?x ∈ {0..1} ∧ path-image ?p = r'{0..?x} ∧ path-image ?q = r'{?x..1}
        using vts-split-path-image[of r vts ?p ?p-vts ?q ?q-vts ?i - ?x]
      by (smt (verit, ccfv-SIG) add commute add-diff-cancel-left' assms(1) assms(3)
atLeastAtMost-iff atLeastatMost-empty' image-empty le-add1 less-diff-conv path-image-nonempty
polygon-of-def)
      thus ?thesis by (metis atLeastAtMost-iff image-Un ivl-disj-un-two-touch(4)
path-image-def)
    qed
    moreover have f'H = convex hull (f'(path-image r))
      using bij by (simp add: calculation(1) convex-hull-linear-image)
    ultimately have H-image: ?H' = f'H by (simp add: image-Un path-image-compose)

  have p-image: path-image ?p' = f'(path-image ?p) using path-image-compose by
blast
  have q-image: path-image ?q' = f'(path-image ?q) using path-image-compose by
blast

  have pathstart-p: pathstart ?p = a
    by (metis Suc-eq-plus1 a-def assms(3) gr-implies-not0 hd-conv-nth length-tl
less-Suc-eq-0-disj list.sel(2) list.size(3) nth-take polygon-pathstart take-eq-Nil)
  have pathfinish-p: pathfinish ?p = b

```

by (*metis* (*no-types*, *lifting*) *H-def H-r add-diff-cancel-right'* *assms*(3) *b-def convex-hull-eq-empty length-take less-add-one less-diff-conv min.absorb4 nth-append one-neq-zero path-image-nonempty polygon-pathfinish set-empty take-eq-Nil vts-split zero-eq-add-iff-both-eq-0*)

then have *pathstart-q*: *pathstart ?q = b using assms*(3) *b-def polygon-pathstart*
by *force*

have *pathstart-p'*: *pathstart ?p' = f a using pathstart-compose pathstart-p* **by**
blast

have *pathfinish-p'*: *pathfinish ?p' = f b using pathfinish-compose pathfinish-p* **by**
blast

have *pathstart-q'*: *pathstart ?q' = f b using pathstart-compose pathstart-q* **by**
blast

have *sublist ?p-vts vts* **by** *auto*

then have *lf-p*: *loop-free ?p*

by (*metis* *add.commute assms*(1) *assms*(3) *less-diff-conv less-imp-le-nat polygon-def polygon-of-def simple-path-def take-i-is-loop-free trans-le-add2*)

then have *simple-p*: *simple-path ?p*
using *assms unfolding polygon-of-def*
by (*meson make-polygonal-path-gives-path simple-path-def*)

have *sublist ?q-vts vts* **by** *auto*

then have *lf-q*: *loop-free ?q*

by (*metis* (*no-types*, *lifting*) *Suc-1 Suc-diff-Suc assms*(1) *assms*(3) *diff-is-0-eq drop-i-is-loop-free less-Suc-eq-le less-zeroE linorder-not-less polygon-def polygon-of-def simple-path-def*)

then have *simple-q*: *simple-path ?q*
using *assms unfolding polygon-of-def*
by (*meson make-polygonal-path-gives-path simple-path-def*)

have *bounded-linear*: *bounded-linear f using bij linear-conv-bounded-linear* **by**
blast

have *1*: *simple-path ?p'*
using *simple-p simple-path-linear-image bij bij-is-inj bounded-linear*
by *blast*

have *2*: *simple-path ?q'*
using *simple-q simple-path-linear-image bij bij-is-inj bounded-linear*
by *blast*

have *3*: *path-image ?p' ∩ path-image ?q' = {pathstart ?q', pathfinish ?q'}*

proof –

have *path-image ?p ∩ path-image ?q ⊆ {pathstart ?q, pathfinish ?q}*
using *loop-free-split-int[of r vts ?p-vts i ?q-vts ?p ?q]*

by (*smt* (*verit*, *ccfv-threshold*) *a-def add-diff-cancel-right'* *assms*(1) *assms*(3) *constant-linepath-is-not-loop-free drop-eq-Nil have-wraparound-vertex hd-conv-nth insert-commute last-conv-nth last-drop last-snoc le-add2 less-diff-conv lf-q linorder-not-less loop-free-split-int make-polygonal-path.simps*(1) *pathstart-p polygon-def polygon-of-def polygon-pathfinish simple-path-def*)

moreover have *pathstart ?q ∈ path-image ?q ∧ pathfinish ?q ∈ path-image ?q*

by *blast*

moreover have $\text{pathstart } ?q \in \text{path-image } ?p \wedge \text{pathfinish } ?q \in \text{path-image } ?p$

by (*smt (verit, ccfv-SIG) a-def add-diff-cancel-right' assms(1) assms(3) b-def constant-linepath-is-not-loop-free drop-eq-Nil have-wraparound-vertex hd-conv-nth last-conv-nth last-drop last-snoc length-take less-add-one less-diff-conv lf-q linorder-not-less list.size(3) make-polygonal-path.simps(1) min.absorb4 nth-take pathfinish-in-path-image pathstart-in-path-image pathstart-p pathstart-q polygon-of-def polygon-pathfinish take-eq-Nil zero-eq-add-iff-both-eq-0 zero-neq-one*)

ultimately have $\text{path-image } ?p \cap \text{path-image } ?q = \{\text{pathstart } ?q, \text{pathfinish } ?q\}$ by *fast*

moreover have $\text{path-image } ?p' \cap \text{path-image } ?q' = f'(\text{path-image } ?p \cap \text{path-image } ?q)$

by (*metis bij bij-is-inj image-Int p-image q-image*)

ultimately show *?thesis* by (*simp add: pathfinish-compose pathstart-compose*)

qed

have 4: $(\text{pathstart } ?p')\$1 = 0 \wedge (\text{pathstart } ?p')\$2 = 0 \wedge (\text{pathfinish } ?p')\$2 = 0$

proof–

have $f ?x = ?e$ using *f* by *blast*

then have $f b - f a = ?e$

by (*metis assms(8) diff-zero f norm-eq-zero orthogonal-transformation-norm*)

moreover have $f a = 0$ by (*metis assms(8) f norm-eq-zero orthogonal-transformation-norm*)

moreover from *calculation* have $f b = ?e$ by *force*

ultimately show *?thesis* using *pathfinish-p' pathstart-p'* by *auto*

qed

have 5: $(\text{pathfinish } ?p')\$1 > 0$

proof–

have $\text{pathfinish } ?p' = f b$ using *pathfinish-p'* by *auto*

moreover have $f b = ?e$ using *assms(8) f* by *auto*

moreover have $?e\$1 = \text{norm } ?x$ by *simp*

ultimately show *?thesis* using *a-neq-b* by *auto*

qed

have 6: $?H' = \text{convex hull } (\text{path-image } ?p' \cup \text{path-image } ?q')$ by *blast*

have 7: $\{\text{pathstart } ?p', \text{pathfinish } ?p'\} \subseteq \text{frontier } ?H'$

proof–

have $\{\text{pathstart } ?p, \text{pathfinish } ?p\} \subseteq \text{frontier } H$

using *pathstart-p pathfinish-p assms(6)* by *fastforce*

then have $f'\{\text{pathstart } ?p, \text{pathfinish } ?p\} \subseteq f'(\text{frontier } H)$ by *blast*

moreover have $f'(\text{frontier } H) = \text{frontier } (f'H)$

by (*simp add: bij bij-is-inj frontier-injective-linear-image*)

ultimately show *?thesis* using *H-image* by (*simp add: pathfinish-compose pathstart-compose*)

qed

have 8: $?p'\{0 < .. < 1\} \subseteq \text{interior } ?H'$

proof–

have 1: *convex H* by (*simp add: H-def*)

have 2: $\forall j \in \{0 < .. < \text{length } ?p\text{-vts} - 1\}. ?p\text{-vts } ! j \notin \text{frontier } H$

by (*simp add: add commute assms(3) assms(7) less-diff-conv*)

have 3: *loop-free ?p* using *lf-p* by *blast*

have 4: $\text{path-image } ?p \subseteq H$ using *H-r hull-subset r-union* by *fastforce*

have 5: $\text{length } ?p\text{-vts} \geq 3$ **using** *assms(3)* **by** *force*
have $?p'\{0 < .. < 1\} \subseteq \text{interior } H$ **using** *vts-interior[OF 1 2 3 4 5]* **by** *argo*
moreover **have** $f'(?p'\{0 < .. < 1\}) = ?p'\{0 < .. < 1\}$ **by** (*meson image-comp*)
moreover **have** $f'(\text{interior } H) = \text{interior } ?H'$
using *H-image interior-injective-linear-image[of f H]* **by** (*simp add: bij*
bij-is-inj)
ultimately show *?thesis* **by** *fast*
qed
have 9: $\text{pathstart } ?q' = \text{pathfinish } ?p' \wedge \text{pathfinish } ?q' = \text{pathstart } ?p'$
by (*metis (mono-tags, lifting) H-def H-r a-def assms(1) constant-linepath-is-not-loop-free*
convex-hull-eq-empty drop-eq-Nil have-wraparound-vertex hd-conv-nth last-conv-nth
last-drop last-snoc lf-q linorder-not-less make-polygonal-path.simps(1) path-image-nonempty
pathfinish-compose pathfinish-p pathstart-compose pathstart-p pathstart-q polygon-of-def
polygon-pathfinish set-empty)

let $?l = \text{open-segment } a \ b$
let $?l' = \text{open-segment } (\text{pathstart } ?p') \ (\text{pathfinish } ?p')$

have *: $\text{path-image } ?q' \cap \text{open-segment } (\text{pathstart } ?p') \ (\text{pathfinish } ?p') = \{\} \wedge$
 $?l' \subseteq \text{frontier } ?H'$
using *pocket-fill-line-int-aux8[OF 1 2 3 4 5 6 7 8 9]* **by** *blast*
moreover **have** *l-image: ?l' = f' ?l*
proof –
have $f a = \text{pathstart } ?p' \wedge f b = \text{pathfinish } ?p'$ **using** *pathfinish-p' pathstart-p'*
by *presburger*
moreover **have** $\bigwedge a \ b. f'(\text{open-segment } a \ b) = \text{open-segment } (f a) \ (f b)$
by (*simp add: bij bij-is-inj open-segment-linear-image*)
ultimately show *?thesis* **by** *presburger*
qed
moreover **have** $\text{path-image } ?q' = f'(\text{path-image } ?q)$ **using** *q-image* **by** *blast*
ultimately **have** $\text{path-image } ?q \cap ?l = \{\}$ **by** *blast*
moreover **have** $\text{path-image } ?p \cap ?l = \{\}$
proof –
from 8 **have** $\text{path-image } ?p' \cap ?l' = \{\}$
proof –
have $?p'\{0 < .. < 1\} \cap ?l' = \{\}$
by (*smt (verit, ccfv-SIG) * 8 Diff-disjoint disjoint-iff frontier-def subset-iff*)
moreover **have** $?p' \ 0 \notin ?l'$
by (*metis * 9 IntI empty-iff pathfinish-in-path-image pathstart-def*)
moreover **have** $?p' \ 1 \notin ?l'$
by (*metis * 9 Int-iff emptyE pathfinish-def pathstart-in-path-image*)
ultimately show *?thesis*
by (*smt (verit, ccfv-SIG) * 1 3 9 Int-Un-eq(4) Un-Diff-cancel Un-iff dis-*
joint-iff insert-commute simple-path-endless)
qed
thus *?thesis* **using** *l-image bij p-image* **by** *auto*
qed
ultimately **have** $\text{path-image } r \cap ?l = \{\}$
by (*simp add: r-union boolean-algebra.conj-disj-distrib inf-commute*)

moreover have $a \in \text{path-image } r$ **using** *pathstart-p r-union* **by** *auto*
moreover have $b \in \text{path-image } r$ **using** *pathfinish-p r-union* **by** *auto*
moreover have $(\text{path-image } (\text{linepath } a \ b)) = ?l \cup \{a, b\}$ **by** (*simp add: closed-segment-eq-open*)
ultimately show $\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } r = \{a, b\}$ **by** *auto*

have $l'\text{-frontier: } ?l \subseteq \text{frontier } ?H'$ **using** $*$ **by** *presburger*
have $?l \subseteq \text{frontier } H$
proof –
have $?l' = f' ?l$ **using** *l-image* **by** *blast*
moreover have $\text{frontier } ?H' = f'(\text{frontier } H)$
by (*metis H-image bij bij-is-inj frontier-injective-linear-image*)
ultimately have $f' ?l \subseteq f'(\text{frontier } H)$ **using** $l'\text{-frontier}$ **by** *argo*
thus $?thesis$ **by** (*simp add: bij bij-is-inj inj-image-subset-iff*)
qed
moreover have $\text{closed-segment } a \ b = \text{path-image } (\text{linepath } a \ b)$ **by** *simp*
moreover have $\text{closed-segment } a \ b = ?l \cup \{a, b\}$ **by** (*simp add: closed-segment-eq-open*)
moreover have $a \in \text{frontier } H \wedge b \in \text{frontier } H$ **using** *assms(6)* **by** *auto*
ultimately show $\text{path-image } (\text{linepath } a \ b) \subseteq \text{frontier } H$ **by** *simp*
qed

lemma *linepath-translation*: $(\lambda v. v - a) \circ (\text{linepath } x \ y) = \text{linepath } ((\lambda v. v - a) \ x) ((\lambda v. v - a) \ y)$
by (*auto simp: linepath-def algebra-simps*)

lemma *linepath-image-translation*:
 $\text{path-image } ((\lambda v. v - a) \circ (\text{linepath } x \ y)) = \text{path-image } (\text{linepath } ((\lambda v. v - a) \ x) ((\lambda v. v - a) \ y))$
using *linepath-translation* **by** *metis*

lemma *make-polygonal-path-translate*:
assumes $\text{length } vts \geq 1$
shows $(\lambda v. v - a) \circ (\text{make-polygonal-path } vts) = \text{make-polygonal-path } (\text{map } (\lambda v. v - a) \ vts)$
using *assms*
proof (*induct length vts arbitrary: vts a*)
case 0
then show $?case$ **by** *linarith*
next
case (*Suc n*)
{ assume $*$: *Suc n = 1*
then have $\text{make-polygonal-path } vts = \text{linepath } (vts!0) \ (vts!0)$
by (*metis Cons-nth-drop-Suc One-nat-def Suc.hyps(2) Suc.prem1 drop0 drop-eq-Nil less-numeral-extra(1) make-polygonal-path.simps(2)*)
then have $(\lambda v. v - a) \circ (\text{make-polygonal-path } vts) = \text{linepath } ((vts!0) - a) \ ((vts!0) - a)$
by *fastforce*
then have $?case$
by (*metis Cons-nth-drop-Suc One-nat-def Suc.hyps(2) Suc.prem1 * drop0*)

```

drop-eq-Nil list.map(1) list.simps(9) make-polygonal-path.simps(2) zero-less-one)
} moreover
{ assume *: Suc n = 2
  then have make-polygonal-path vts = linepath (vts!0) (vts!1)
  by (metis (no-types, lifting) Cons-nth-drop-Suc One-nat-def Suc.hyps(2) Suc-1
diff-Suc-1 drop0 drop-Suc drop-eq-Nil le-numeral-extra(4) length-tl less-numeral-extra(1)
make-polygonal-path.simps(3) nth-tl pos2)
  then have  $(\lambda v. v - a) \circ (\text{make-polygonal-path } vts) = \text{linepath } ((vts!0) - a)
((vts!1) - a)$ 
  using linepath-translation by auto
  then have ?case
  by (metis (no-types, lifting) * Cons-nth-drop-Suc One-nat-def Suc.hyps(2)
Suc-1 drop0 drop-eq-Nil length-map lessI make-polygonal-path.simps(3) nat-le-linear
nth-map pos2)
} moreover
{ assume *: Suc n  $\geq$  3
  then obtain h h' t where vts: vts = h # h' # t
  by (metis Suc.hyps(2) Suc-le-length-iff numeral-3-eq-3)
  then have  $(\lambda v. v - a) \circ (\text{make-polygonal-path } (h' \# t))
= \text{make-polygonal-path } (\text{map } (\lambda v. v - a) (h' \# t))$ 
  using Suc.hyps(1) Suc.hyps(2) * by auto
  moreover have  $(\lambda v. v - a) \circ (\text{linepath } h h') = \text{linepath } (h - a) (h' - a)$ 
  using linepath-translation by blast
  moreover have make-polygonal-path vts = (linepath h h') +++ (make-polygonal-path
(h' # t))
  by (metis * Suc.hyps(2) Suc-le-length-iff vts list.sel(3) make-polygonal-path.simps(4)
numeral-3-eq-3)
  ultimately have ?case
  by (smt (verit) list.discI list.inject list.simps(9) make-polygonal-path.elims
path-compose-join vts)
}
ultimately show ?case using Suc.premis by linarith
qed

```

lemma pocket-fill-line-int:

```

assumes polygon-of r vts
defines H  $\equiv$  convex hull (set vts)
assumes  $2 \leq i \wedge i < \text{length } vts - 1$ 
defines a  $\equiv$  hd vts
defines b  $\equiv$  vts!i
assumes  $\{a, b\} \subseteq \text{frontier } H$ 
assumes  $\forall j \in \{0 <..<i\}. vts!j \notin \text{frontier } H$ 
shows path-image (linepath a b)  $\cap$  path-image r =  $\{a, b\}$ 
  path-image (linepath a b)  $\subseteq$  frontier H

```

proof –

```

let ?f =  $(\lambda v. v - a)::(\text{real}^2 \Rightarrow \text{real}^2)$ 
let ?r' = ?f  $\circ$  r
let ?vts' = map ?f vts
let ?H' = convex hull (set ?vts')

```

```

let ?a' = ?f a
let ?b' = ?f b

have 5: hd ?vts' = 0
by (metis One-nat-def a-def assms(3) cancel-comm-monoid-add-class.diff-cancel
lessI list.map-sel(1) list.size(3) nat-diff-split-asm not-less-zero)

have a'b': ?a' = hd ?vts' ∧ ?b' = ?vts'!i using 5 assms(3) b-def by force

have frontier-H': frontier ?H' = ?f '(frontier H)
using frontier-translation[of -a H]
by (metis (no-types, lifting) H-def convex-hull-translation image-cong list.set-map
uminus-add-conv-diff)

have simple-path r using assms(1) polygon-def polygon-of-def by blast
then have simple-path ?r' using simple-path-translation-eq[of -a r] by simp
moreover have ?r' = make-polygonal-path ?vts'
using make-polygonal-path-translate assms(1) assms(3) polygon-of-def by auto
moreover have closed-path ?r'
by (smt (verit, best) closed-path-def add-diff-inverse-nat assms(1) assms(3) cal-
culation(1) calculation(2) dual-order.refl gr-implies-not0 hd-conv-nth length-map
less-Suc-eq-le list.map-disc-iff list.map-sel(1) nat-diff-split-asm nth-map plus-1-eq-Suc
polygon-def polygon-of-def polygon-pathfinish polygon-pathstart simple-path-def)
ultimately have 1: polygon-of ?r' ?vts'
unfolding polygon-of-def polygon-def polygon-def polygonal-path-def by blast
have 2: 2 ≤ i ∧ i < length ?vts' - 1 using assms(3) by auto
have 3: {hd ?vts', ?vts'!i} ⊆ frontier ?H'
using a'b' frontier-H'
by (metis (no-types, lifting) assms(6) image-empty image-insert image-mono)
have 4: ∀j ∈ {0 <..<i}. ?vts'!j ∉ frontier ?H'
proof
fix j assume *: j ∈ {0 <..<i}
then have vts!j ∉ frontier H using assms(7) by blast
then have ?f (vts!j) ∉ frontier ?H' using frontier-H' by auto
thus ?vts'!j ∉ frontier ?H' using Nat.le-imp-diff-is-add * assms(3) by auto
qed

have path-image (linepath ?a' ?b') ∩ path-image ?r' = {?a', ?b'}
using pocket-fill-line-int-0(1)[OF 1 2 3 4 5] a'b' by argo
moreover have {?a', ?b'} = ?f'{a, b} by simp
moreover have path-image (linepath ?a' ?b') = ?f'(path-image (linepath a b))
using linepath-image-translation path-image-compose by blast
moreover have path-image ?r' = ?f'(path-image r) using path-image-compose
by blast
ultimately have ?f'(path-image (linepath a b)) ∩ ?f'(path-image r) = ?f'{a, b}
by argo
then have ?f'(path-image (linepath a b) ∩ path-image r) = ?f'{a, b} by (simp
add: image-Int)
moreover have bij ?f by (simp add: bij-diff-right)

```

ultimately show $\text{path-image } (\text{linepath } a \ b) \cap \text{path-image } r = \{a, b\}$
by (*meson bij-is-inj inj-image-eq-iff*)

have $\text{path-image } (\text{linepath } ?a' \ ?b') \subseteq \text{frontier } ?H'$
using *pocket-fill-line-int-0(2)[OF 1 2 3 4 5] a'b' by argo*
thus $\text{path-image } (\text{linepath } a \ b) \subseteq \text{frontier } H$
by (*metis <bij ?f> <path-image (linepath ?a' ?b') = ?f'(path-image (linepath a b))> bij-betw-imp-inj-on frontier-H' inj-image-subset-iff*)
qed

lemma *path-connected-simple-path-endless*:
assumes *simple-path p*
shows *path-connected (path-image p - {pathstart p, pathfinish p}) (is path-connected ?S)*
proof –
have *continuous-on {0<..
using *assms(1) unfolding simple-path-def path-def*
by (*meson continuous-on-path dual-order.refl greaterThanLessThan-subseteq-atLeastAtMost-iff path-def*)
moreover have *path-connected {0<..
ultimately have *path-connected (p'{0<..
by *blast*
thus *?thesis using simple-path-endless assms by metis*
qed***

lemma *simple-loop-split*:
assumes *simple-path p ∧ closed-path p*
assumes *simple-path q*
assumes $\text{path-image } q \cap \text{path-image } p = \{q \ 0, q \ 1\}$
assumes $\text{path-image } q \cap \text{path-inside } p \neq \{\}$
shows $q\{0<..
proof –
have *inside-outside: inside-outside p (path-inside p) (path-outside p)*
using *Jordan-inside-outside-real2 closed-path-def assms(1) inside-outside-def path-inside-def path-outside-def*
by *presburger*$

obtain *x where x: x ∈ path-image q ∩ path-inside p using assms(4) by blast*
then obtain *tx where tx ∈ {0..1} ∧ q tx = x unfolding path-image-def by fast*
moreover then have $tx \neq 0 \wedge tx \neq 1$
using *assms(3) inside-outside x unfolding inside-outside-def by auto*
ultimately have *tx: tx ∈ {0<..*

have *connected (q'{0<..
using *connected-simple-path-endless simple-path-endless assms(2) by metis*
then have *path-connected (q'{0<..
using *path-connected-simple-path-endless assms(2) simple-path-endless by metis***

moreover have $q^{\{0 < .. < 1\}} \cap \text{path-inside } p \neq \{\}$ **using** *tx x* **by** *blast*
moreover have $q^{\{0 < .. < 1\}} \cap \text{frontier } (\text{path-inside } p) = \{\}$
using *inside-outside unfolding inside-outside-def*
by (*smt (verit, del-insts) Diff-Int-distrib2 assms(2,3) diff-eq inf-compl-bot-right inf-idem inf-sup-aci(1) pathfinish-def pathstart-def simple-path-endless*)
ultimately show *?thesis*
using *path-connected-not-frontier-subset*[of $q^{\{0 < .. < 1\}}$ *path-inside* *p*] **by** *fast*
qed

lemma *pocket-path-interior-aux:*

assumes *simple-path* $p \wedge$ *simple-path* q
assumes *arc* $p \wedge$ *arc* q
assumes $q\ 0 = p\ 1 \wedge q\ 1 = p\ 0$
assumes $\text{path-image } p \cap \text{path-image } q = \{p\ 0, q\ 0\}$
defines $A \equiv \text{convex hull } (\text{path-image } p \cup \text{path-image } q)$
defines $l \equiv \text{linepath } (p\ 0) (p\ 1)$
assumes $p^{\{0 < .. < 1\}} \subseteq \text{interior } A$
assumes $\text{path-image } l \subseteq \text{frontier } A$
assumes $\text{path-image } q \cap \text{path-image } l = \{l\ 0, q\ 0\}$
shows $p^{\{0 < .. < 1\}} \cap \text{path-inside } (l\ +++\ q) \neq \{\}$
 $\text{simple-path } (l\ +++\ q) \wedge \text{closed-path } (l\ +++\ q)$
 $\text{path-image } p \cap \text{path-image } (l\ +++\ q) = \{p\ 0, p\ 1\}$

proof –

let $?r = l\ +++\ q$
let $?Ir = \text{path-inside } ?r$
let $?Or = \text{path-outside } ?r$
show *closed-simple-r: simple-path* $?r \wedge$ *closed-path* $?r$
using *simple-path-join-loop*[of $l\ q$] *assms unfolding pathstart-def pathfinish-def*
by (*metis (no-types, opaque-lifting) closed-path-def arc-linepath arc-simple-path dual-order.refl inf-commute linepath-0' linepath-1' pathfinish-def pathfinish-join pathstart-def pathstart-join simple-path-def*)
then have *inside-outside-r: inside-outside* $?r\ ?Ir\ ?Or$
by (*simp add: Jordan-inside-outside-real2 closed-path-def inside-outside-def path-inside-def path-outside-def*)

have *l-p-endpoints: l 0 = p 0 \wedge l 1 = p 1* **by** (*simp add: l-def linepath-0' linepath-1'*)

have *l-q-endpoints: l 0 = q 1 \wedge l 1 = q 0* **by** (*simp add: assms(3) l-p-endpoints*)

have *p-int-l: p* $^{\{0 < .. < 1\}}$ *\cap path-image* $l = \{\}$ **using** *assms(7,8) unfolding frontier-def* **by** *blast*

have *q-int-l: q* $^{\{0 < .. < 1\}}$ *\cap path-image* $l = \{\}$

by (*metis (no-types, opaque-lifting) assms(9) Diff-iff Int-Diff all-not-in-conv assms(1) assms(3) inf-sup-aci(1) insert-commute l-def linepath-0' pathfinish-def pathstart-def simple-path-endless*)

have *interval: $\{0..1::\text{real}\} = \{0 < .. < 1\} \cup \{0, 1\}$* **by** *fastforce*

have *lf-l: loop-free* l

using *closed-simple-r not-loop-free-first-component simple-path-def* **by** *blast*

let $?p' = \text{reversepath } p$

```

let ?s = l +++ ?p'
let ?Is = path-inside ?s
let ?Os = path-outside ?s
have arc ?p'  $\wedge$  arc l
  by (metis assms(2) arc-linepath arc-reversepath arc-simple-path l-def pathfin-
ish-def pathstart-def)
moreover have p'-int-l: path-image ?p'  $\cap$  path-image l = {?p' 0, l 0}
proof–
  have path-image p  $\cap$  path-image l = {l 0, l 1}
  proof–
    have {l 0, l 1}  $\subseteq$  path-image p  $\cap$  path-image l
      using assms(3) assms(4) l-def linepath-0' linepath-1' by fastforce
    moreover have path-image p = p'{0<..\cup {p 0, p 1}
      using interval unfolding path-image-def by blast
    ultimately show ?thesis using p-int-l l-p-endpoints by simp
  qed
  moreover have ?p' 0 = l 1 by (simp add: l-def linepath-1' reversepath-def)
  moreover have path-image p = path-image ?p' by simp
  ultimately show ?thesis by (metis doubleton-eq-iff)
qed
ultimately have closed-simple-s: closed-path ?s  $\wedge$  simple-path ?s
  using simple-path-join-loop[of l ?p'] assms unfolding pathstart-def pathfin-
ish-def
  by (metis (no-types, opaque-lifting) closed-path-def dual-order.refl inf-commute
insert-commute linepath-0' linepath-1' pathfinish-def pathfinish-join pathfinish-reversepath
pathstart-def pathstart-join pathstart-reversepath simple-path-def)
  then have inside-outside-s: inside-outside ?s ?Is ?Os
  by (simp add: Jordan-inside-outside-real2 closed-path-def inside-outside-def
path-inside-def path-outside-def)

  have r-inside-subset: path-inside ?r  $\subseteq$  interior A
  proof–
    have path-image l  $\subseteq$  A  $\wedge$  path-image q  $\subseteq$  A
    by (metis A-def Un-upper2 assms(1) assms(8) compact-Un compact-convex-hull
compact-simple-path-image frontier-subset-compact hull-subset subset-trans)
    thus ?thesis
    by (metis (no-types, lifting) A-def closed-simple-r convex-contains-simple-closed-path-imp-contains-path-ins
convex-convex-hull inside-outside-def inside-outside-r interior-eq interior-mono sub-
set-path-image-join)
  qed
  have s-inside-subset: path-inside ?s  $\subseteq$  interior A
  proof–
    have path-image l  $\subseteq$  A  $\wedge$  path-image p  $\subseteq$  A
    by (metis A-def Un-upper1 assms(1) assms(8) compact-Un compact-convex-hull
compact-simple-path-image frontier-subset-compact hull-subset subset-trans)
    thus ?thesis
    by (metis A-def Jordan-inside-outside-real2 closed-path-def closed-simple-s
convex-contains-simple-closed-path-imp-contains-path-inside convex-convex-hull in-
terior-maximal path-image-reversepath path-inside-def subset-path-image-join)

```

qed

have $q\{0 < .. < 1\} \subseteq \text{path-outside } ?s$
proof(rule ccontr)
 let $?ep = \{v. v \text{ extreme-point-of } A\}$
 assume $\neg q\{0 < .. < 1\} \subseteq \text{path-outside } ?s$
 then have $\exists x \in q\{0 < .. < 1\}. x \in \text{path-inside } ?s \cup \text{path-image } ?s$
 using *inside-outside-s unfolding inside-outside-def* **by** *auto*
 then have $q\{0 < .. < 1\} \subseteq \text{path-inside } ?s$
 using *simple-loop-split*[of $p \ q$]
 by (smt (verit) *DiffE IntI Int-Un-distrib2 closed-path-def UnE \arc (reversepath p) \wedge \arc l \succ \arc-imp-path assms(1) assms(2) assms(3) assms(4) closed-simple-r closed-simple-s doubleton-eq-iff emptyE inf.commute l-def path-image-join path-image-reversepath path-join-eq pathfinish-join pathfinish-linepath pathstart-join pathstart-linepath simple-loop-split simple-path-endless simple-path-joinE sup-absorb2*)
 then have $q\{0 < .. < 1\} \cap \text{frontier } A = \{\}$ **using** *frontier-def s-inside-subset* **by** *fastforce*
 then have $(\text{path-image } p \cup \text{path-image } q) \cap \text{frontier } A = \{p \ 0, \ p \ 1\}$
 by (smt (z3) *Diff-disjoint Int-Un-distrib Un-Diff-Int Un-Int-eq(3) assms(1) assms(3) assms(4) assms(7) assms(8) assms(9) frontier-def inf.commute inf.orderE inf-idem inf-left-commute insert-commute l-p-endpoints pathfinish-def pathstart-def simple-path-endless*)
 moreover have $?ep \subseteq \text{path-image } p \cup \text{path-image } q$
 by (*simp add: extreme-points-of-convex-hull A-def*)
 moreover have $?ep \subseteq \text{frontier } A$
 using *extreme-point-not-in-interior*
 proof–
 have $?ep \cap \text{interior } A = \{\}$
 using *extreme-point-not-in-interior* **by** *blast*
 thus *?thesis*
 by (smt (verit, ccfv-SIG) *A-def Int-Un-distrib2 Un-Diff-cancel assms(1) calculation(2) closure-convex-hull compact-Un compact-simple-path-image dual-order.trans frontier-def hull-subset inf.absorb-iff2 inf-commute sup-bot-left*)
 qed
 ultimately have $*$: $?ep \subseteq \{p \ 0, \ p \ 1\}$ **by** *auto*
 have $A = \text{path-image } l$
 proof–
 have $\text{convex } A \wedge \text{compact } A$
 by (*simp add: A-def arc-imp-path assms(2) compact-Un compact-convex-hull compact-path-image*)
 then have $A\text{-ep}: A = \text{convex hull } ?ep$ **using** *Krein-Milman-Minkowski* **by** *blast*
 moreover have *finite ?ep* **using** $*$ *infinite-super* **by** *auto*
 moreover have $A \neq \{\}$ **by** (*simp add: A-def*)
 moreover have $\forall x. A \neq \{x\}$ **using** *assms(7)* **by** *fastforce*
 ultimately have $\text{card } ?ep \geq 2$ **using** *convex-hull-two-extreme-points* **by** *metis*
 then have $?ep = \{p \ 0, \ p \ 1\}$
 by (*metis * One-nat-def Suc-1 add-leD2 card.empty card-insert-disjoint card-seteq finite.emptyI finite.insertI insert-absorb plus-1-eq-Suc*)

then have $A = \text{closed-segment } (p \ 0) \ (p \ 1)$ **by** $(\text{metis } A\text{-ep } \text{segment-convex-hull})$
thus $?thesis$ **by** $(\text{simp add: } l\text{-def})$
qed
then have $\text{interior } A = \{\}$
by $(\text{metis } A\text{-def } \text{Diff-eq-empty-iff } \text{assms}(1) \ \text{assms}(8) \ \text{closure-convex-hull}$
 $\text{compact-Un } \text{compact-simple-path-image } \text{double-diff } \text{dual-order.refl } \text{frontier-def } \text{interior-subset})$
thus False **using** $\text{inside-outside-def } \text{inside-outside-r } \text{r-inside-subset}$ **by** auto
qed

let $?e = l \ (1/2)$
have $l\text{-on-r-frontier: } \text{path-image } l \subseteq \text{frontier } (\text{path-inside } ?r)$
using $\text{inside-outside-r } \text{unfolding } \text{inside-outside-def}$
by $(\text{metis } \text{Un-upper1 } \text{closed-simple-r } \langle \text{arc } (\text{reversepath } p) \rangle \wedge \text{arc } l \rangle \text{arc-def}$
 $\text{assms}(2) \ \text{path-image-join } \text{path-join-eq } \text{simple-path-def})$
moreover have $\text{path-image } l \subseteq \text{frontier } (\text{path-inside } ?s)$
using $\text{inside-outside-s } \text{unfolding } \text{inside-outside-def}$
by $(\text{simp add: } l\text{-def } \text{path-image-join } \text{pathstart-def } \text{reversepath-def})$
ultimately have $e\text{-frontier: } ?e \in \text{frontier } (\text{path-inside } ?r) \wedge ?e \in \text{frontier}$
 $(\text{path-inside } ?s)$
by $(\text{simp add: } \text{path-defs}(4) \ \text{subsetD})$

have $e\text{-notin: } ?e \notin \text{path-image } p \cup \text{path-image } q$
proof–
have $?e \notin \text{path-image } p$
proof–
have $?e \neq l \ 0 \wedge ?e \neq l \ 1$ **using** $lf\text{-l } \text{unfolding } \text{loop-free-def}$ **by** fastforce
then have $?e \neq p \ 0 \wedge ?e \neq p \ 1$ **using** $l\text{-p-endpoints}$ **by** simp
moreover have $?e \notin p\{0 < .. < 1\}$ **using** $p\text{-int-l } \text{unfolding } \text{path-image-def}$
by fastforce
ultimately show $?thesis$ **using** $p\text{-int-l } \text{unfolding } \text{path-image-def}$ **by** fastforce
qed
moreover have $?e \notin \text{path-image } q$
proof–
have $?e \neq l \ 0 \wedge ?e \neq l \ 1$ **using** $lf\text{-l } \text{unfolding } \text{loop-free-def}$ **by** fastforce
then have $?e \neq q \ 0 \wedge ?e \neq q \ 1$ **using** $l\text{-q-endpoints}$ **by** simp
moreover have $?e \notin q\{0 < .. < 1\}$ **using** $q\text{-int-l } \text{unfolding } \text{path-image-def}$
by fastforce
ultimately show $?thesis$ **using** $q\text{-int-l } \text{unfolding } \text{path-image-def}$ **by** fastforce
qed
ultimately show $?thesis$ **by** blast
qed

obtain ε **where** $\varepsilon: \varepsilon > 0 \wedge \text{ball } ?e \ \varepsilon \cap \text{path-image } p = \{\} \wedge \text{ball } ?e \ \varepsilon \cap \text{path-image}$
 $q = \{\}$
proof–
have $?e \notin \text{path-image } p$ **using** $e\text{-notin}$ **by** simp
moreover have $\text{compact } (\text{path-image } p)$ **by** $(\text{simp add: } \text{assms}(2) \ \text{compact-arc-image})$
moreover have $?e \notin \text{path-image } q$ **using** $e\text{-notin}$ **by** simp
moreover have $\text{compact } (\text{path-image } q)$ **by** $(\text{simp add: } \text{assms}(2) \ \text{compact-arc-image})$

ultimately obtain $\varepsilon 1 \ \varepsilon 2$ **where**
 $\varepsilon 1 > 0 \wedge \text{ball } ?e \ \varepsilon 1 \cap \text{path-image } p = \{\} \wedge \varepsilon 2 > 0 \wedge \text{ball } ?e \ \varepsilon 2 \cap \text{path-image } q = \{\}$
by (*meson* *assms*(1) *not-on-path-ball* *simple-path-imp-path*)
thus *?thesis* **using** *that*[of *min* $\varepsilon 1 \ \varepsilon 2$] **by** (*simp* *add*: *disjoint-iff*)
qed

obtain *z-r* **where** *z-r*: $z-r \in \text{ball } ?e \ \varepsilon \cap \text{path-inside } ?r$
by (*metis* *e-frontier* ε *all-not-in-conv* *disjoint-iff* *frontier-straddle* *mem-ball*)
obtain *z-s* **where** *z-s*: $z-s \in \text{ball } ?e \ \varepsilon \cap \text{path-inside } ?s$
by (*metis* *e-frontier* ε *all-not-in-conv* *disjoint-iff* *frontier-straddle* *mem-ball*)

have *z-s-in-r*: $z-s \in \text{path-inside } ?r$
proof–
let *?l-z* = *linepath* *z-r* *z-s*
have *z-r* \in *interior* *A* \wedge *z-s* \in *interior* *A*
using *r-inside-subset* *s-inside-subset* *z-r* *z-s* **by** *blast*
then **have** *path-image* *?l-z* \subseteq *interior* *A* **by** (*simp* *add*: *A-def* *closed-segment-subset*)
then **have** *1*: *path-image* *?l-z* \cap *path-image* *l* = $\{\}$
by (*smt* (*verit*) *Diff-iff* *assms*(8) *disjoint-iff* *frontier-def* *subsetD*)

have *convex* (*ball* *?e* ε) **by** *simp*
then **have** *path-image* *?l-z* \subseteq *ball* *?e* ε
by (*metis* *IntD1* *closed-segment-subset* *path-image-linepath* *z-r* *z-s*)
then **have** *2*: *path-image* *?l-z* \cap *path-image* *q* = $\{\}$ **using** ε **by** *blast*

show *?thesis*
by (*smt* (*verit*, *best*) *1* *2* *IntI* *Int-Un-distrib* *Int-Un-distrib2* *Jordan-inside-outside-real2* *closed-path-def* ε $\langle \text{path-image } (\text{linepath } z-r \ z-s) \subseteq \text{ball } (l \ (1 / 2)) \ \varepsilon \rangle$ *arc-def* *assms*(2) *closed-simple-r* *emptyE* *in-mono* *inf.assoc* *le-iff-inf* *path-connected-not-frontier-subset* *path-connected-path-image* *path-image-join* *path-inside-def* *path-join-path-ends* *path-linepath* *pathfinish-in-path-image* *pathfinish-linepath* *pathstart-in-path-image* *pathstart-linepath* *sup.order-iff* *z-r*)
qed

let *?xq* = *q* (1/2)
let *?z* = *z-s*

let *?v* = *?xq* – *?z*
let *?ray* = $\lambda d. ?z + d *_{\mathbb{R}} ?v$
let *?rayline* = *linepath* *?z* *?xq*
have *z-ray*: *?z* = *?ray* 0 **by** *simp*
have *xq-ray*: *?xq* = *?ray* 1 **by** *simp*
have *xq-rayline*: *?xq* = *?rayline* 1 **unfolding** *linepath-def* **by** *simp*
have *?xq* \in *path-image* *?r*
by (*metis* (*mono-tags*, *opaque-lifting*) *Un-iff* *atLeastAtMost-iff* *imageI* *l-q-endpoints* *less-eq-real-def* *path-defs*(4) *path-image-join* *pathfinish-def* *pathstart-def* *pos-half-less* *zero-less-divide-1-iff* *zero-less-numeral* *zero-less-one*)
then **have** *xq-frontier*: *?xq* \in *frontier* (*path-inside* *?r*)

```

    using inside-outside-r unfolding inside-outside-def by auto
  have xq-neq-z: ?xq ≠ ?z
  proof -
    have ?xq ∈ path-image ?r
    proof -
      have q (1 / 2) ∈ path-image q
      by (simp add: path-defs(4))
      thus ?thesis
      by (simp add: l-q-endpoints path-image-join pathfinish-def pathstart-def)
    qed
    thus ?thesis using z-s-in-r inside-outside-r unfolding inside-outside-def by
blast
  qed
  then have v-neq-0: ?v ≠ 0 by simp

  have bounded (path-inside ?r) using inside-outside-r unfolding inside-outside-def
  by blast
  moreover have ?z ∈ interior (path-inside ?r)
  by (metis inside-outside-def inside-outside-r interior-eq z-s-in-r)
  ultimately obtain d where d: 0 < d ∧ ?ray d ∈ frontier (path-inside ?r)
  ∧ (∀ e ∈ {0..<d}. ?ray e ∈ interior (path-inside ?r))
  using ray-to-frontier[of path-inside ?r ?z ?v] by (metis atLeastLessThan-iff
v-neq-0)

  have interior-inside-r: interior (path-inside ?r) = path-inside ?r
  by (meson inside-outside-def inside-outside-r interior-eq)
  have d-leq-1: d ≤ 1
  proof (rule ccontr)
    assume ¬ d ≤ 1
    then have d > 1 by simp
    moreover have ?ray 1 ∈ frontier (path-inside ?r) using xq-ray xq-frontier by
  argo
  ultimately show False using d unfolding frontier-def by fastforce
  qed

  have z-inside: ?z ∈ path-inside ?s using z-s by blast
  moreover have ?rayline d ∈ path-outside ?s
  proof -
    have ?rayline d ∉ path-image l if d < 1
    proof -
      have ?rayline 0 ∈ interior A
      using r-inside-subset by (simp add: linepath-0' subsetD z-s-in-r)
      moreover have path-image ?rayline ⊆ closure A
      proof -
        have closure A = A
        using A-def assms(1) closure-convex-hull compact-Un compact-simple-path-image
      by blast
      moreover have ?rayline 0 ∈ A using ⟨?rayline 0 ∈ interior A⟩ inte-
rior-subset by blast

```

```

moreover have ?rayline 1 ∈ A
  using path-image-def A-def hull-subset xq-rayline by fastforce
ultimately show ?thesis
  by (metis A-def closed-segment-subset convex-convex-hull linepath-0'
linepath-1' path-image-linepath)
qed
moreover have ¬ path-image ?rayline ⊆ rel-frontier A
proof–
  have path-image ?rayline ∩ interior A ≠ {}
    using ⟨?rayline 0 ∈ interior A⟩ unfolding path-image-def by fastforce
  moreover have interior A ∩ rel-frontier A = {}
    using rel-frontier-def rel-interior-nonempty-interior by auto
  ultimately show ?thesis by blast
qed
ultimately have rel-interior (path-image ?rayline) ⊆ rel-interior A
  using subset-rel-interior-convex[of path-image ?rayline A] by (simp add:
A-def)
moreover have interior A = rel-interior A
  using ⟨?rayline 0 ∈ interior A⟩ rel-interior-nonempty-interior by auto
moreover have ?rayline d ∈ ?rayline{0<..<1} using that d by simp
ultimately show ?thesis
  by (smt (verit, del-Insts) DiffD1 DiffD2 Un-iff xq-neq-z arc-linepath arc-simple-path
assms(8) closed-segment-eq-open frontier-def path-image-linepath pathfinish-linepath
pathstart-linepath rel-interior-closed-segment simple-path-endless subset-eq)
qed
moreover have ?rayline d ∉ path-image l if d = 1
  using that q-int-l unfolding linepath-def by (simp add: disjoint-iff)
moreover have ?rayline d ∈ path-image ?r
  by (metis (no-types, lifting) add-diff-eq d diff-add-eq inside-outside-def in-
side-outside-r linepath-def scale-left-diff-distrib scale-one scale-right-diff-distrib)
ultimately show ?thesis
  by (smt (verit, ccfv-SIG) d-leq-1 Diff-iff Int-iff closed-path-def ⟨arc (reversepath
p) ∧ arc l⟩ arc-def assms(1) assms(3) assms(9) closed-simple-r insert-commute
l-def l-p-endpoints not-in-path-image-join path-join-eq pathfinish-join pathfinish-linepath
pathstart-join pathstart-linepath q-outside simple-path-def simple-path-endless sub-
setD)
qed
moreover have ?z ∈ ?rayline{0..d}
  using z-ray unfolding linepath-def
  by (smt (verit, del-Insts) add commute atLeastAtMost-iff cancel-comm-monoid-add-class.diff-cancel
d diff-zero image-iff less-eq-real-def segment-degen-1)
moreover have ?rayline d ∈ ?rayline{0..d} by (simp add: d less-eq-real-def)
ultimately have ?rayline{0..d} ∩ path-inside ?s ≠ {} ∧ ?rayline{0..d} ∩
path-outside ?s ≠ {}
  by blast
then have ?rayline{0..d} ∩ path-inside ?s ≠ {} ∧ ?rayline{0..d} ∩ ¬ path-inside
?s ≠ {}
  using inside-outside-s unfolding inside-outside-def by (meson ComplI dis-
joint-iff)

```

moreover have $\text{path-connected } (?rayline\{0..d\})$
proof–
have $?rayline\{0..d\} = \text{path-image } (\text{subpath } 0\ d\ ?rayline)$ **by** (*simp add: d path-image-subpath*)
moreover have $\text{path } (\text{subpath } 0\ d\ ?rayline)$ **using** $d\ d\text{-leq-1}$ **by** *auto*
ultimately show $?thesis$ **by** (*metis path-connected-path-image*)
qed
ultimately have $?rayline\{0..d\} \cap \text{frontier } (\text{path-inside } ?s) \neq \{\}$
using $\text{path-connected-frontier}[of\ ?rayline\{0..d\}\ \text{path-inside } ?s]$ **by** (*metis disjoint-iff*)
then have $?rayline\{0..d\} \cap \text{path-image } ?s \neq \{\}$ **using** *inside-outside-s unfolding inside-outside-def* **by** *argo*
moreover have $?rayline\ 0 \notin \text{path-image } ?s$
proof–
have $?xq \neq p\ 0$
by (*metis (full-types) disjoint-iff greaterThanLessThan-iff imageI l-p-endpoints pathstart-def pathstart-in-path-image pos-half-less q-int-l zero-less-divide-1-iff zero-less-numeral zero-less-one*)
moreover have $?xq \neq p\ 1$
by (*metis (full-types) disjoint-iff greaterThanLessThan-iff imageI l-p-endpoints pathfinish-def pathfinish-in-path-image pos-half-less q-int-l zero-less-divide-1-iff zero-less-numeral zero-less-one*)
moreover have $?xq \notin p\{0<..
proof–
have $?xq \in q\{0<.. **by** *fastforce*
thus $?thesis$ **by** (*metis assms(1,3,4) Diff-iff Int-iff pathfinish-def pathstart-def simple-path-endless*)
qed
moreover have $?xq \notin \text{path-image } l$
by (*metis disjoint-iff greaterThanLessThan-iff imageI pos-half-less q-int-l zero-less-divide-1-iff zero-less-numeral zero-less-one*)
ultimately show $?thesis$
by (*metis (no-types, lifting) ComplD UnI1 z-inside inside-outside-def inside-outside-s linopath-0'*)
qed
moreover have $?rayline\ d \notin \text{path-image } ?s$
using $\langle ?rayline\ d \in \text{path-outside } ?s \rangle$ *inside-outside-def inside-outside-s* **by** *auto*
moreover have $\{0..d\} = \{0<.. **using** d **by** *fastforce*
ultimately have $?rayline\{0<.. **unfolding** *path-image-def*
by *blast*
moreover have $?rayline\{0<..
unfolding *linopath-def* **by** (*auto simp: algebra-simps*)
moreover have $?ray\{0<.. **using** d *interior-inside-r* **by** *fastforce*
ultimately have $\text{path-image } ?s \cap \text{path-inside } ?r \neq \{\}$ **by** *blast*
moreover have $\text{path-image } l \cap \text{path-inside } ?r = \{\}$
by (*metis (no-types, opaque-lifting) Diff-disjoint Int-assoc l-on-r-frontier frontier-def inf.orderE inf-bot-left inf-sup-aci(1) interior-inside-r*)
moreover have $p\{0<..$$$$$$$

proof–

have $path\text{-}image\ ?s = path\text{-}image\ p \cup path\text{-}image\ l$
by (*simp add: l-p-endpoints path-image-join pathfinish-def sup-commute*)
moreover have $p\{0<..
by (*metis assms(1) pathfinish-def pathstart-def simple-path-endless*)
ultimately have $path\text{-}image\ ?s = p\{0<..
using *assms(3) assms(9) l-p-endpoints* **by** *auto*
moreover have $p\ 1 \in path\text{-}image\ l \wedge p\ 0 \in path\text{-}image\ l$ **by** (*simp add: l-def*)
ultimately show *?thesis* **using** *p-int-l* **by** *blast*$$

qed

ultimately show $p\{0<.. **by** *auto*$

show $path\text{-}image\ p \cap path\text{-}image\ (l\ +++\ q) = \{p\ 0, p\ 1\}$

by (*smt (verit, best) Int-Un-distrib Un-absorb assms(1) assms(3) assms(4)*
closed-simple-r insert-commute l-p-endpoints p'-int-l path-image-join path-image-reversepath
path-join-path-ends reversepath-def simple-path-imp-path)

qed

lemma *pocket-path-interior*:

assumes *simple-path p* \wedge *simple-path q*
assumes *arc p* \wedge *arc q*
assumes $q\ 0 = p\ 1 \wedge q\ 1 = p\ 0$
assumes $path\text{-}image\ p \cap path\text{-}image\ q = \{p\ 0, q\ 0\}$
defines $A \equiv convex\ hull\ (path\text{-}image\ p \cup path\text{-}image\ q)$
defines $l \equiv linepath\ (p\ 0)\ (p\ 1)$
assumes $p\{0<..
assumes $path\text{-}image\ l \subseteq frontier\ A$
assumes $path\text{-}image\ q \cap path\text{-}image\ l = \{l\ 0, q\ 0\}$
shows $p\{0<..
using *pocket-path-interior-aux*[*of p q*] *simple-loop-split*[*of l +++ q p*] *assms*
by (*metis (no-types, lifting) DiffE disjoint-iff simple-path-endless*)$$

lemma *pocket-path-good*:

assumes *polygon* (*make-polygonal-path vts*)
assumes $vts!0 \in frontier\ (convex\ hull\ (set\ vts))$
assumes $vts!1 \notin frontier\ (convex\ hull\ (set\ vts))$
assumes $\neg convex\ (path\text{-}image\ (make\text{-}polygonal\text{-}path\ vts) \cup path\text{-}inside\ (make\text{-}polygonal\text{-}path\ vts))$
defines $pocket\text{-}path\text{-}vts \equiv construct\text{-}pocket\text{-}0\ vts\ (set\ vts \cap frontier\ (convex\ hull\ (set\ vts)))$
defines $pocket \equiv make\text{-}polygonal\text{-}path\ (pocket\text{-}path\text{-}vts\ @\ [pocket\text{-}path\text{-}vts!0])$
defines $filled\text{-}vts \equiv fill\text{-}pocket\text{-}0\ vts\ (length\ pocket\text{-}path\text{-}vts)$
defines $filled\text{-}p \equiv make\text{-}polygonal\text{-}path\ filled\text{-}vts$
defines $a \equiv hd\ pocket\text{-}path\text{-}vts$
defines $b \equiv last\ pocket\text{-}path\text{-}vts$
defines $good\text{-}pocket\text{-}path\text{-}vts \equiv tl\ (butlast\ pocket\text{-}path\text{-}vts)$
shows *polygon filled-p*
 $is\text{-}polygon\text{-}split\text{-}path\ (butlast\ filled\text{-}vts)\ 0\ 1\ good\text{-}pocket\text{-}path\text{-}vts$
polygon pocket

```

      card (set pocket-path-vts) < card (set vts)
      card (set filled-vts) < card (set vts)
proof –
  let ?p = make-polygonal-path vts
  let ?A = set vts ∩ frontier (convex hull (set vts))
  let ?filled-vts-tl = tl filled-vts
  let ?filled-p-tl = make-polygonal-path ?filled-vts-tl
  let ?pocket-vts = pocket-path-vts @ [pocket-path-vts!0]
  let ?pocket-path = make-polygonal-path pocket-path-vts
  let ?l = linepath a b

  let ?r = min-nonzero-index-in-set vts ?A
  have int-A-nonempty: set (tl vts) ∩ ?A ≠ {}
  by (metis (mono-tags, lifting) IntI Nitpick.size-list-simp(2) Suc-eq-plus1 assms(1)
  assms(2) card-length empty-iff have-wraparound-vertex last-in-set last-tl le-add1
  le-trans not-less-eq-eq numeral-3-eq-3 polygon-at-least-3-vertices snoc-eq-iff-butlast)
  then have r-defined: nonzero-index-in-set vts ?A ?r ∧ (∀ i < ?r. ¬ nonzero-index-in-set
  vts ?A i)
  using min-nonzero-index-in-set-defined[of vts ?A] by fast

  have two-vts-on-frontier: 2 ≤ card ?A
  by (metis convex-hull-two-vts-on-frontier One-nat-def Suc-1 add-leD2 assms(1)
  numeral-3-eq-3 plus-1-eq-Suc polygon-at-least-3-vertices)
  moreover have frontier-vts-subset: ?A ⊆ set vts by force
  moreover have distinct-vts: distinct (butlast vts)
  using assms(1) polygon-def simple-polygonal-path-vts-distinct by blast
  moreover have hd-last-vts: hd vts = last vts
  by (metis assms(1) have-wraparound-vertex hd-conv-nth snoc-eq-iff-butlast)
  ultimately have a-neq-b: a ≠ b
  using a-def b-def construct-pocket-0-first-last-distinct pocket-path-vts-def by
  presburger
  have length filled-vts ≥ 2
  unfolding filled-vts-def fill-pocket-0-def
  by (smt (verit, best) One-nat-def Suc-1 Suc-diff-Suc a-def a-neq-b b-def con-
  struct-pocket-0-def diff-is-0-eq diff-zero hd-Nil-eq-last length-drop length-greater-0-conv
  length-tl list.sel(3) not-less-eq-eq pocket-path-vts-def sublist-length-le sublist-take)
  moreover have filled-vts-0: a = filled-vts!0
  unfolding filled-vts-def fill-pocket-0-def a-def pocket-path-vts-def construct-pocket-0-def
  by auto
  moreover have filled-vts-1: b = filled-vts!1
  by (smt (verit, del-insts) filled-vts-def fill-pocket-0-def b-def pocket-path-vts-def
  construct-pocket-0-def Cons-nth-drop-Suc Nitpick.size-list-simp(2) a-def a-neq-b add.right-neutral
  drop0 drop-eq-Nil hd-Nil-eq-last last-conv-nth length-take length-tl linorder-not-less
  list.sel(3) min.absorb4 nat-le-linear not-less-eq-eq nth-drop nth-take plus-1-eq-Suc
  take-all-iff zero-less-diff)
  ultimately have filled-vts: filled-vts = [a, b] @ tl ?filled-vts-tl
  by (metis (no-types, lifting) Nitpick.size-list-simp(2) One-nat-def Suc-1 ap-

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pend-Nil append-eq-Cons-conv length-greater-0-conv list.collapse not-less-eq-eq nth-Cons-0 nth-tl order-less-le-trans pos2)

have 1: *polygon-of ?p vts unfolding polygon-of-def using assms(1) by blast*
have 2: $2 \leq ?r \wedge ?r < \text{length } vts - 1$
proof–
have $?r \neq 0 \wedge ?r \neq 1$
using *assms(2,3) min-nonzero-index-in-set-def nonzero-index-in-set-def r-defined*
by *fastforce*
then have 1: $?r \geq 2$ **by** *simp*

have $\exists i \in \{0 <..< \text{length } vts - 1\}. vts!i \in \text{frontier } (\text{convex hull } (\text{set } vts))$
proof–
have $\text{card } ((\text{set } vts) \cap \text{frontier } (\text{convex hull } (\text{set } vts))) \geq 2$
using *two-vts-on-frontier by blast*
then obtain *v* **where** $v \in \text{set } vts \wedge v \in \text{frontier } (\text{convex hull } \text{set } vts) \wedge v \neq$
hd vts
by (*metis hd-last-vts Int-iff a-neq-b assms(2) b-def construct-pocket-0-last-in-set convex-hull-empty empty-set fill-pocket-0-def filled-vts-0 filled-vts-def frontier-empty hd-conv-nth int-A-nonempty last-in-set nth-Cons-0 pocket-path-vts-def*)
thus *?thesis*
by (*metis hd-last-vts assms(1) in-set-conv-nth diff-Suc-1 gr0-implies-Suc greaterThanLessThan-iff have-wraparound-vertex last-conv-nth le-eq-less-or-eq less-Suc-eq-le less-one nat.simps(3) nat-le-linear snoc-eq-iff-butlast*)
qed
then have 2: $?r < \text{length } vts - 1$
using *r-defined*
unfolding *min-nonzero-index-in-set-def nonzero-index-in-set-def*
by (*smt (verit, del-insts) Int-iff add commute add-diff-cancel-left' add-diff-inverse-nat greaterThanLessThan-iff less-imp-diff-less mem-Collect-eq nat-less-le nth-mem*)
show *?thesis using 1 2 by blast*
qed
have *ab*: $a = \text{hd } vts \wedge b = vts! ?r$
by (*metis (no-types, lifting) 2 Suc-1 int-A-nonempty ab-semigroup-add-class.add-ac(1) add-Suc-right b-def construct-pocket-0-def fill-pocket-0-def filled-vts-0 filled-vts-def hd-drop-conv-nth last-snoc le-add-diff-inverse2 min-nonzero-index-in-set-bound nth-Cons-0 plus-1-eq-Suc pocket-path-vts-def take-hd-drop*)
have 3: $\{\text{hd } vts, vts ! ?r\} \subseteq \text{frontier } (\text{convex hull } \text{set } vts)$
using *ab assms(1) assms(2) assms(3) b-def construct-pocket-is-pocket is-pocket-0-def pocket-path-vts-def*
by *fastforce*
have 4: $\forall j \in \{0 <..< ?r\}. vts ! j \notin \text{frontier } (\text{convex hull } \text{set } vts)$
using *r-defined unfolding nonzero-index-in-set-def by fastforce*

have *l-int-p*: $\text{path-image } (\text{linepath } (\text{hd } vts) (vts ! ?r)) \cap \text{path-image } ?p = \{\text{hd } vts, vts ! ?r\}$
using *pocket-fill-line-int[OF 1 2 3 4] by blast*
have *l-frontier*: $\text{path-image } (\text{linepath } (\text{hd } vts) (vts ! ?r)) \subseteq \text{frontier } (\text{convex hull } (\text{set } vts))$

```

using pocket-fill-line-int[OF 1 2 3 4] by blast

have path-image ?filled-p-tl  $\cap$  path-image ?l = {a, b}
proof –
  have path-image (linepath (hd vts) (vts ! ?r))  $\cap$  path-image ?p = {hd vts, vts !
  ?r}
    using pocket-fill-line-int[OF 1 2 3 4] by blast
  moreover have path-image ?filled-p-tl  $\subseteq$  path-image ?p
  proof –
    have sublist ?filled-vts-tl vts by (simp add: fill-pocket-0-def filled-vts-def)
    thus ?thesis using  $\langle 2 \leq \text{length filled-vts} \rangle$  sublist-path-image-subset by auto
  qed
  moreover have a  $\in$  path-image ?filled-p-tl  $\wedge$  b  $\in$  path-image ?filled-p-tl
    by (smt (verit, best) Cons-nth-drop-Suc Diff-insert-absorb One-nat-def Suc-1
     $\langle 2 \leq \text{length filled-vts} \rangle$  drop0 drop-eq-Nil fill-pocket-0-def filled-vts-0 filled-vts-1 filled-vts-def
    hd-last-vts last-drop last-in-set linorder-not-le list.sel(3) not-less-eq-eq nth-Cons-0
    order-less-le-trans pathstart-in-path-image polygon-pathstart pos2 subset-Diff-insert
    vertices-on-path-image)
  ultimately show ?thesis using ab by auto
  qed
  moreover have hd-filled: hd ?filled-vts-tl = last [a, b]
    unfolding filled-vts-def fill-pocket-0-def pocket-path-vts-def construct-pocket-0-def
    by (metis construct-pocket-0-def fill-pocket-0-def filled-vts filled-vts-def hd-append2
    last-ConsL last-ConsR list.sel(1) list.sel(3) list.simps(3) pocket-path-vts-def tl-append2)
  moreover have last-filled: last ?filled-vts-tl = hd [a, b]
    unfolding filled-vts-def fill-pocket-0-def pocket-path-vts-def construct-pocket-0-def
    using r-defined a-def assms(1) assms(2) assms(3) construct-pocket-is-pocket
    hd-last-vts is-pocket-0-def pocket-path-vts-def
    by fastforce
  moreover have loop-free ?filled-p-tl
  proof –
    have sublist ?filled-vts-tl vts
    unfolding filled-vts-def fill-pocket-0-def pocket-path-vts-def construct-pocket-0-def
    using r-defined
    by force
    thus ?thesis
    by (smt (verit, del-insts) Nitpick.size-list-simp(2) Suc-1  $\langle 2 \leq \text{length filled-vts} \rangle$ 
     $\langle b = \text{filled-vts} ! 1 \rangle$  a-neq-b assms(1) diff-is-0-eq dual-order.strict-trans1 last-conv-nth
    last-filled le-antisym length-greater-0-conv length-tl list.sel(1) list.size(3) not-less-eq-eq
    nth-tl polygon-def pos2 simple-path-def sublist-is-loop-free sublist-length-le)
  qed
  moreover have loop-free ?l using a-neq-b linepath-loop-free by blast
  moreover have filled-vts: filled-vts = [a, b] @ tl ?filled-vts-tl using filled-vts by
  blast
  moreover have arc ?l
    by (smt (verit) arc-linepath calculation(5) constant-linepath-is-not-loop-free)
  moreover have arc ?filled-p-tl
    by (smt (z3) arc-simple-path calculation(2) calculation(3) calculation(4) cal-

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culation(7) hd-Nil-eq-last hd-conv-nth last.simps last-conv-nth list.discI list.sel(1)
make-polygonal-path-gives-path pathfinish-linepath pathstart-linepath polygon-pathfinish
polygon-pathstart simple-path-def)
moreover have $?l = \text{make-polygonal-path } [a, b]$
using *make-polygonal-path.simps* **by** *presburger*
ultimately have *lf-filled: loop-free filled-p*
by (*smt (z3) Nat.add-diff-assoc One-nat-def Suc-pred' add-Suc-shift append-butlast-last-id*
arc-distinct-ends butlast.simps(2) filled-p-def hd-Nil-eq-last hd-conv-nth inf-sup-aci(1)
last-ConsR less-numeral-extra(1) list.sel(1) list.simps(3) list.size(3) list.size(4)
loop-free-append nth-append-length order-eq-refl plus-1-eq-Suc polygon-pathfinish poly-
gon-pathstart)
show *polygon-filled-p: polygon filled-p*
unfolding *polygon-def*
by (*metis closed-path-def UNIV-def append-is-Nil-conv filled-p-def filled-vts*
hd-append2 last.simps last-conv-nth last-filled lf-filled list.discI list.exhaust-sel make-polygonal-path-gives-path
nth-Cons-0 polygon-pathfinish polygon-pathstart polygonal-path-def rangeI simple-path-def)

have $\{a, b\} \subseteq \text{set filled-vts}$
using *filled-vts* **by** (*smt (z3) UnCI empty-set list.simps(15) set-append sub-*
set-iff)
moreover have *pocket-path: ?pocket-path = make-polygonal-path ([a] @ good-pocket-path-vts*
@ [b])
by (*metis (no-types, lifting) a-def a-neq-b append-Cons append-Nil append-butlast-last-id*
b-def good-pocket-path-vts-def hd-Nil-eq-last hd-conv-nth last-conv-nth length-butlast
list.collapse list.size(3) tl-append2)
moreover have *path-image ?pocket-path \subseteq path-inside filled-p $\cup \{a, b\}$*
proof –
let $?p = ?\text{pocket-path}$
let $?q = ?\text{filled-p-tl}$
let $?H = \text{convex hull } (\text{path-image } ?p \cup \text{path-image } ?q)$
have $b: \text{pocket-path-vts} = \text{take } (?r + 1) \text{ vts}$
unfolding *pocket-path-vts-def construct-pocket-0-def* **by** *blast*
moreover then have $c': ?\text{filled-vts-tl} = \text{drop } ?r \text{ vts}$ **unfolding** *filled-vts-def*
fill-pocket-0-def
using *2* **by** *fastforce*
ultimately have $\text{vts} = \text{pocket-path-vts} @ \text{tl } ?\text{filled-vts-tl}$
by (*metis Suc-eq-plus1 append-take-drop-id drop-Suc tl-drop)*
then have *path-image ?p = path-image ?p \cup path-image ?q*
by (*metis Suc-1 a-def a-neq-b b-def diff-is-0-eq hd-Nil-eq-last hd-conv-nth*
hd-filled last.simps last-conv-nth last-filled list.discI list.sel(1) make-polygonal-path-image-append-alt
not-less-eq-eq path-image-join polygon-pathfinish polygon-pathstart)
moreover have *convex hull (path-image ?p) = convex hull (set vts)*
by (*metis (no-types, lifting) 1 Un-subset-iff convex-hull-of-polygon-is-convex-hull-of-vts*
hull-Un-subset hull-mono subset-antisym vertices-on-path-image)
ultimately have $H\text{-eq: } ?H = \text{convex hull } (\text{set vts})$ **by** *presburger*

have $a: ?p = \text{make-polygonal-path } \text{vts} \wedge \text{loop-free } ?p$
using *assms(1) polygon-def simple-path-def* **by** *blast*

have c : $?filled\text{-}vts\text{-}tl = drop ((?r + 1) - 1) vts$ **using** c' **by** *simp*
have h : $1 \leq ?r + 1 \wedge ?r + 1 < length\ vts$ **using** 2 **by** *linarith*
have $path\text{-}image\ ?p \cap path\text{-}image\ ?q \subseteq \{?p\ 0, ?q\ 0\}$
using *loop-free-split-int*[*OF a b c - - - h*] **by** (*simp add: pathstart-def*)
moreover have $?p\ 0 \in path\text{-}image\ ?p \wedge ?p\ 0 \in path\text{-}image\ ?q$
by (*metis a-def a-neq-b b-def hd-Nil-eq-last hd-conv-nth hd-filled last.simps last-conv-nth last-filled list.sel(1) pathfinish-in-path-image pathstart-def pathstart-in-path-image polygon-pathfinish polygon-pathstart*)
moreover have $?q\ 0 \in path\text{-}image\ ?p \wedge ?q\ 0 \in path\text{-}image\ ?q$
by (*metis a-def a-neq-b b-def hd-Nil-eq-last hd-conv-nth hd-filled last.simps last-conv-nth last-filled list.sel(1) pathfinish-in-path-image pathstart-def pathstart-in-path-image polygon-pathfinish polygon-pathstart*)
ultimately have 4 : $path\text{-}image\ ?p \cap path\text{-}image\ ?q = \{?p\ 0, ?q\ 0\}$ **by** *fastforce*

have 1 : $simple\text{-}path\ ?p \wedge simple\text{-}path\ ?q$
by (*metis (no-types, lifting) One-nat-def Suc-1 Suc-le-eq <arc ?filled-p-tl> arc-simple-path assms(1) assms(2) assms(3) construct-pocket-is-pocket is-pocket-0-def le-add2 make-polygonal-path-gives-path numeral-3-eq-3 order-le-less-trans plus-1-eq-Suc pocket-path-vts-def polygon-def simple-path-def sublist-is-loop-free sublist-take*)
have 2 : $arc\ ?p \wedge arc\ ?q$
by (*metis 1 <arc ?filled-p-tl> a-def a-neq-b b-def hd-Nil-eq-last hd-conv-nth last-conv-nth polygon-pathfinish polygon-pathstart simple-path-cases*)
have 3 : $?q\ 0 = ?p\ 1 \wedge ?q\ 1 = ?p\ 0$
by (*metis 1 a-def append-Cons b-def constant-linepath-is-not-loop-free filled-vts hd-conv-nth last-conv-nth last-filled list.sel(1) list.sel(3) make-polygonal-path.simps(1) pathfinish-def pathstart-def polygon-pathfinish polygon-pathstart simple-path-def*)
have 5 : $?p\ \{0 <.. < 1\} \subseteq interior\ ?H$
proof –
have $\forall j \in \{0 <.. < ?r\}. vts!j \notin frontier\ (convex\ hull\ (set\ vts))$
by (*smt (verit, del-insts) Int-iff dual-order.strict-trans greaterThanLessThan-iff int-A-nonempty mem-Collect-eq min-nonzero-index-in-set-defined nonzero-index-in-set-def nth-mem*)
moreover have $?r = length\ pocket\text{-}path\text{-}vts - 1$ **using** $b\ h$ **by** *auto*
moreover have $\forall j < ?r. vts!j = pocket\text{-}path\text{-}vts!j$ **using** b **by** *auto*
ultimately have $\forall j \in \{0 <.. < length\ pocket\text{-}path\text{-}vts - 1\}. pocket\text{-}path\text{-}vts!j \notin frontier\ ?H$
using $H\text{-eq}$ **by** *simp*
moreover have *loop-free ?pocket-path* **using** 1 *simple-path-def* **by** *auto*
ultimately show $?thesis$
by (*metis vts-interior Un-subset-iff assms(1) assms(2) assms(3) construct-pocket-is-pocket convex-convex-hull hull-subset is-pocket-0-def pocket-path-vts-def*)
qed
have 6 : $path\text{-}image\ (linepath\ (?p\ 0)\ (?p\ 1)) \subseteq frontier\ ?H$
by (*metis l-frontier H-eq 3 a-def a-neq-b ab b-def hd-Nil-eq-last hd-conv-nth hd-filled last.simps last-filled list.discI list.sel(1) pathstart-def polygon-pathstart*)
have 7 : $path\text{-}image\ ?q \cap path\text{-}image\ (linepath\ (?p\ 0)\ (?p\ 1)) = \{linepath\ (?p\ 0)\ (?p\ 1)\ 0, ?q\ 0\}$
by (*metis 3 <path-image (make-polygonal-path (tl filled-vts)) > &cap path-image (linepath a b) = \{a, b\}> a-def a-neq-b b-def hd-Nil-eq-last hd-filled last.simps last-conv-nth*)

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last-filled linepath-0' list.sel(1) pathfinish-def polygon-pathfinish)
  have ?p ' {0<..

```

have $01: 0 < \text{length (butlast filled-vts)} \wedge 1 < \text{length (butlast filled-vts)}$
by (*metis One-nat-def Suc-lessI filled-vts-1 filled-vts-as-butlast a-neq-b append-eq-Cons-conv filled-0-a length-greater-0-conv nth-Cons-Suc nth-append-length*)
show *is-split-path*:
is-polygon-split-path (butlast filled-vts) 0 1 good-pocket-path-vts
using *good-polygonal-path-implies-polygon-split-path*
[OF polygon-filled-p filled-p-as-butlast - 01 filled-0-a filled-1-b le]
using *good-polygonal-path filled-vts-as-butlast*
by *presburger*

have *polygon-pocket-rev: polygon (make-polygonal-path (a#([] @ [b] @ (rev good-pocket-path-vts) @ [a])))*
unfolding *is-polygon-split-path-def*
by (*smt (z3) 01 One-nat-def add-diff-cancel-left' add-diff-cancel-right' filled-0-a filled-1-b is-polygon-split-path-def is-split-path nth-butlast plus-1-eq-Suc take0*)
moreover have *rev-pocket-vts: rev ?pocket-vts = a#([] @ [b] @ (rev good-pocket-path-vts) @ [a])*
by (*smt (verit) a-def a-neq-b append.left-neutral append-Cons append-butlast-last-id b-def good-pocket-path-vts-def hd-Nil-eq-last hd-append2 hd-conv-nth last-conv-nth length-butlast list.collapse list.size(3) rev.simps(1) rev.simps(2) rev-append*)
ultimately show *polygon pocket*
by (*metis polygon-pocket-rev rev-vts-is-polygon polygon-of-def pocket-def rev-rev-ident*)

have $\text{card (set vts)} = \text{length (butlast vts)}$
using *distinct-vts*
by (*smt (verit, ccfv-threshold) Suc-n-not-le-n Un-insert-right append-Nil2 assms(1) butlast-conv-take distinct-card dual-order.strict-trans have-wraparound-vertex hd-conv-nth hd-in-set hd-take insert-absorb length-0-conv length-butlast less-eq-Suc-le linorder-linear list.set(2) not-numeral-le-zero numeral-3-eq-3 polygon-at-least-3-vertices-wraparound polygon-vertices-length-at-least-4 set-append*)
then have $\text{set pocket-path-vts} \subset \text{set vts}$
unfolding *pocket-path-vts-def construct-pocket-0-def*
using *r-defined*
by (*smt (verit, ccfv-threshold) Cons-nth-drop-Suc One-nat-def Suc-diff-Suc Suc-le-lessD add-diff-cancel-right' assms(1) assms(2) assms(3) butlast-conv-take butlast-snoc card-length construct-pocket-0-def construct-pocket-is-pocket drop0 fill-pocket-0-def filled-vts-def is-pocket-0-def is-polygon-split-path-def is-split-path leD le-less-Suc-eq length-butlast length-drop length-greater-0-conv list.inject numeral-3-eq-3 plus-1-eq-Suc pocket-path-vts-def polygon-at-least-3-vertices-wraparound psubsetI set-take-subset take-eq-Nil add-eq-0-iff-both-eq-0 add-gr-0 cancel-comm-monoid-add-class.diff-cancel diff-zero dual-order.strict-trans filled-p-def length-Cons length-tl less-imp-diff-less list.sel(3) list.size(3) not-less-eq-eq polygon-filled-p zero-less-one zero-neq-one*)
thus $\text{card (set pocket-path-vts)} < \text{card (set vts)}$ **by** (*simp add: psubset-card-mono*)

have $\text{card (set vts)} = \text{card (set (butlast vts))}$
by (*smt (z3) Cons-nth-drop-Suc List.finite-set One-nat-def Suc-1 Suc-le-lessD*)

two-vts-on-frontier distinct-vts hd-last-vts frontier-vts-subset butlast.simps(1) but-last-conv-take card-insert-if card-length card-mono distinct-card drop0 drop-eq-Nil dual-order.trans last-in-set last-tl length-butlast length-greater-0-conv length-tl list.collapse list.sel(3) list.simps(15) set-take-subset verit-la-disequality)

moreover have *length good-pocket-path-vts ≥ 1*
unfolding *good-pocket-path-vts-def pocket-path-vts-def construct-pocket-0-def*
using *convex-hull-of-nonconvex-polygon-strict-subset[OF - assms(4), of vts]*
using *Suc-le-eq assms(1) assms(2) assms(3) construct-pocket-0-def construct-pocket-is-pocket is-pocket-0-def numeral-3-eq-3*
by auto
ultimately show *card (set filled-vts) < card (set vts)*

unfolding *filled-vts-def fill-pocket-0-def good-pocket-path-vts-def pocket-path-vts-def*
by *(smt (verit) Nitpick.size-list-simp(2) Suc-1 Suc-diff-Suc Suc-n-not-le-n <2 \leq length filled-vts> distinct-vts hd-last-vts card-length diff-is-0-eq diff-less distinct-card drop-eq-Nil fill-pocket-0-def filled-vts-def insert-absorb last-drop last-in-set le leI le-less-Suc-eq length-Cons length-butlast length-drop length-tl less-imp-diff-less list.simps(15) order-less-le-trans pocket-path-vts-def)*
qed

29.3 Arbitrary Polygon Case

lemma *pick-rotate:*

assumes *polygon-of p vts*
assumes *all-integral vts*
obtains *p' vts' where polygon-of p' vts'*
 \wedge *vts'!0 \in frontier (convex hull (set vts'))*
 \wedge *path-image p' = path-image p*
 \wedge *all-integral vts'*
 \wedge *set vts' = set vts*
proof –
obtain *v where v: v \in set vts \cap frontier (convex hull (set vts))*
proof –
obtain *v where v \in set vts \wedge v extreme-point-of (convex hull (set vts))*
using *assms unfolding polygon-of-def*
by *(metis List.finite-set card.empty convex-convex-hull convex-hull-eq-empty extreme-point-exists-convex extreme-point-of-convex-hull finite-imp-compact-convex-hull not-numeral-le-zero polygon-at-least-3-vertices)*
then have *v \in set vts \wedge v \in frontier (convex hull (set vts))*
by *(metis Krein-Milman-frontier List.finite-set convex-convex-hull extreme-point-of-convex-hull finite-imp-compact-convex-hull)*
thus *?thesis using that by blast*
qed
obtain *i where i: vts!i = v \wedge i < length vts* **by** *(meson IntE in-set-conv-nth v)*
let *?vts-rotated = rotate-polygon-vertices vts i*
let *?p-rotated = make-polygonal-path ?vts-rotated*
have same-set: *set vts = set ?vts-rotated*
using *assms unfolding polygon-of-def*
using *rotate-polygon-vertices-same-set*

by force
 moreover have *: ?vts-rotated!0 ∈ frontier (convex hull (set ?vts-rotated))
 proof –
 have ?vts-rotated!0 = vts!i
 using assms unfolding polygon-of-def
 by (metis add-leD2 diff-self-eq-0 have-wraparound-vertex hd-conv-nth i last-snoc
 less-nat-zero-code list.size(3) nat-le-linear numeral-Bit0 polygon-vertices-length-at-least-4
 rotated-polygon-vertices)
 moreover have vts!i ∈ frontier (convex hull (set vts)) using v i by blast
 ultimately show ?thesis using same-set by argo
 qed
 moreover have polygon ?p-rotated
 using rotation-is-polygon assms unfolding polygon-of-def by blast
 moreover have all-integral ?vts-rotated
 using rotate-polygon-vertices-same-set assms
 unfolding all-integral-def polygon-of-def by blast
 moreover have path-image ?p-rotated = path-image p
 using assms unfolding polygon-of-def using polygon-vts-arb-rotation by force
 moreover then have path-inside ?p-rotated = path-inside p unfolding path-inside-def
 by simp
 ultimately show ?thesis using polygon-of-def that by blast
 qed

lemma pick-unrotated:

fixes p :: R-to-R2
 assumes polygon: polygon p
 assumes polygonal-path: p = make-polygonal-path vts
 assumes int-vertices: all-integral vts
 assumes I-is: I = card {x. integral-vec x ∧ x ∈ path-inside p}
 assumes B-is: B = card {x. integral-vec x ∧ x ∈ path-image p}
 assumes vts!0 ∈ frontier (convex hull (set vts))
 shows measure lebesgue (path-inside p) = I + B/2 - 1
 using assms
 proof (induct card (set vts) arbitrary: vts p I B rule: less-induct)
 case less
 have B-finite: finite {x. integral-vec x ∧ x ∈ path-image p}
 using finite-path-image less(2) by auto
 have set vts ⊆ {x. integral-vec x ∧ x ∈ path-image p}
 using less(3) vertices-on-path-image[of vts] less(4)
 unfolding all-integral-def
 by auto
 then have card-vts: card (set vts) ≥ 3
 using polygon-at-least-3-vertices[OF less(2) less(3)] card-mono order-trans
 by blast
 have vts-wraparound: vts ! 0 = vts ! (length vts - 1)
 using less(2-3) polygon-pathstart polygon-pathfinish
 unfolding polygon-def closed-path-def
 by (metis diff-0-eq-0 length-0-conv)
 then have vts-is: vts = (butlast vts) @ [vts ! 0]
 end

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by (metis butlast-conv-take have-wraparound-vertex less.prem(1) less.prem(2))
have same-set: set vts = set (butlast (vts))
by (metis ListMem-iff Un-insert-right append.right-neutral butlast.simps(2) constant-linepath-is-not-loop-free elem hd-conv-nth insert-absorb less.prem(1) less.prem(2) list.collapse list.simps(15) make-polygonal-path.simps(2) polygon-def set-append simple-path-def vts-is)
have distinct-butlast-vts: distinct (butlast vts)
  using simple-polygonal-path-vts-distinct less(2-3)
  unfolding polygon-def
by auto
have card-butlast-vts: card (set vts) = card (set (butlast vts))
  using vts-wraparound
  by (smt (verit, best) List.finite-set butlast-conv-take card-distinct card-length card-mono card-vts diff-is-0-eq diff-less distinct-butlast-vts distinct-card drop-rev dual-order.strict-trans1 le-SucE length-append-singleton length-greater-0-conv less-numeral-extra(1) less-numeral-extra(4) nth-eq-iff-index-eq one-less-numeral-iff order-class.order-eq-iff semiring-norm(77) set-drop-subset set-rev vts-is)
then have card-set-len-butlast: card (set vts) = length (butlast vts)
  using distinct-butlast-vts
  by (metis distinct-card)
{ assume triangle: card (set vts) = 3
  then have length (butlast vts) = 3
    using card-set-len-butlast
    by auto
  then have butlast vts = [vts ! 0, vts ! 1, vts ! 2]
    by (metis (no-types, lifting) Cons-nth-drop-Suc One-nat-def Suc-1 card-set-len-butlast card-vts drop0 drop-eq-Nil lessI nth-append numeral-3-eq-3 one-less-numeral-iff semiring-norm(77) vts-is zero-less-numeral)
    then have vts-is: vts = [vts ! 0, vts ! 1, vts ! 2, vts ! 0]
      using vts-is by auto
    then have p-make-triangle: p = make-triangle (vts ! 0) (vts ! 1) (vts ! 2)
      using less(3) unfolding make-triangle-def by simp
    then have not-collinear: ¬ collinear {vts ! 0, vts ! 1, vts ! 2}
      using vts-is less(2) polygon-vts-not-collinear[of p vts] unfolding polygon-of-def make-triangle-def
      by (smt (verit, ccfv-threshold) insert-absorb2 insert-commute list.set(1) list.simps(15))
    have all-integral: all-integral [vts ! 0, vts ! 1, vts ! 2]
      using less.prem(3) vts-is unfolding all-integral-def
      by (simp add: ⟨butlast vts = [vts ! 0, vts ! 1, vts ! 2]⟩ in-set-butlastD)
    have distinct: distinct [vts ! 0, vts ! 1, vts ! 2]
      using ⟨butlast vts = [vts ! 0, vts ! 1, vts ! 2]⟩ distinct-butlast-vts by presburger
    have pick-triangle: pick-triangle p (vts ! 0) (vts ! 1) (vts ! 2)
      using pick-triangle p-make-triangle less(2) not-collinear all-integral distinct
      by simp
    then have ?case
      using pick-triangle-lemma[OF p-make-triangle all-integral distinct not-collinear]
      less.prem(4-5)
      by blast

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} moreover
{ assume non-triangle: card (set vts) > 3
  { assume convex: convex (path-image p ∪ path-inside p)
    then obtain a b where good-linepath a b vts
      using convex-polygon-has-good-linepath non-triangle
      by (metis inf-sup-aci(5) less.prem(1) less.prem(2))
    then have ab-prop: a ≠ b ∧ {a, b} ⊆ set vts ∧ path-image (linepath a b) ⊆
path-inside p ∪ {a, b}
      unfolding good-linepath-def less.prem(2) by presburger
    then have ab-prop-rewrite: a ≠ b ∧ a ∈ set (butlast vts) ∧ b ∈ set (butlast
vts)
      using same-set
      by simp
    have good-linepath-ab: good-linepath a b ((butlast vts) @ [(butlast vts) ! 0])
      using ab-prop vts-is unfolding good-linepath-def
      using ab-prop-rewrite empty-set hd-append2 hd-conv-nth insert-absorb in-
sert-not-empty less.prem(2) same-set
      by (smt (z3))
    then have good-linepath-ba: good-linepath b a ((butlast vts) @ [(butlast vts) !
0])
      using good-linepath-comm good-linepath-def by blast
    obtain i1 j1 where ij-prop: i1 < length (butlast vts) ∧ j1 < length (butlast
vts) ∧
      butlast vts ! i1 = a ∧
      butlast vts ! j1 = b ∧ i1 ≠ j1
      using ab-prop-rewrite
      by (metis distinct-Ex1 distinct-butlast-vts)
    have i-lt-then: i1 < j1 ⇒ is-polygon-split (butlast vts) i1 j1
      using good-linepath-implies-polygon-split[OF less(2), of butlast vts] vts-is
same-set
      using ij-prop good-linepath-ab good-linepath-ba
      by (metis ab-prop-rewrite length-pos-if-in-set less.prem(2) nth-butlast)
    have j-lt-then: j1 < i1 ⇒ is-polygon-split (butlast vts) j1 i1
      using good-linepath-implies-polygon-split[OF less(2), of butlast vts] vts-is
same-set
      using ij-prop good-linepath-ab good-linepath-ba
      by (metis ab-prop-rewrite length-pos-if-in-set less.prem(2) nth-butlast)
    obtain i j where polygon-split: is-polygon-split (butlast vts) i j
      using i-lt-then j-lt-then ij-prop
      by (meson nat-neq-iff)
    then have ij-prop: i < length (butlast vts) ∧ j < length (butlast vts) ∧ i < j
      unfolding is-polygon-split-def
      by blast

have p-is: p = make-polygonal-path (butlast vts @ [butlast vts ! 0])
  using less(3) vts-is
  by (metis length-greater-0-conv nth-butlast same-set set-empty)

```

```

let ?vts1 = take i (butlast vts)
let ?vts2 = take (j - i - 1) (drop (Suc i) (butlast vts))
let ?vts3 = drop (j - i) (drop (Suc i) (butlast vts))

let ?vtsp1 = (butlast vts ! i # ?vts2 @ [butlast vts ! j, butlast vts ! i])
have finite-butlast: finite (set (butlast vts))
  by blast
have vtsp1-subset: set ?vtsp1  $\subseteq$  set (butlast vts)
  using ij-prop
by (smt (verit, del-insts) Un-commute append-Cons append-Nil dual-order.trans
insert-subset list.simps(15) nth-mem set-append set-drop-subset set-take-subset)

let ?p1 = make-polygonal-path ?vtsp1
let ?I1 = card {x. integral-vec x  $\wedge$  x  $\in$  path-inside ?p1}
let ?B1 = card {x. integral-vec x  $\wedge$  x  $\in$  path-image ?p1}
have polygon-p1: polygon ?p1
  using polygon-split unfolding is-polygon-split-def by metis

let ?vtsp2 = ?vts1 @ [butlast vts ! i, butlast vts ! j] @ ?vts3 @ [butlast vts ! 0]
let ?p2 = make-polygonal-path ?vtsp2
have polygon-p2: polygon ?p2
  using polygon-split unfolding is-polygon-split-def by metis

have j-neq: j  $\neq$  i + 1
by (smt (verit, ccfv-SIG) One-nat-def Suc-n-not-le-n Suc-numeral add-Suc-shift
add-implies-diff cancel-ab-semigroup-add-class.diff-right-commute length-Cons length-append
list.size(3) numeral-3-eq-3 plus-1-eq-Suc polygon-p1 polygon-vertices-length-at-least-4
semiring-norm(2) semiring-norm(8) take-eq-Nil)
have subset1: set (take i (butlast vts))  $\subseteq$  set (butlast vts)
  using ij-prop by (meson set-take-subset)
have subset2: set ([butlast vts ! i, butlast vts ! j])  $\subseteq$  set (butlast vts)
  using ij-prop by simp
have subset3: set (take i (butlast vts) @
[butlast vts ! i, butlast vts ! j])  $\subseteq$  set (butlast vts)
  using subset1 subset2 by auto
have subset4: set (drop (j - i) (drop (Suc i) (butlast vts)) @ [butlast vts ! 0])
 $\subseteq$  set (butlast vts)
  using ij-prop set-drop-subset
  by (metis (no-types, opaque-lifting) Un-commute append-Cons append-Nil
card-set-len-butlast drop0 drop-drop drop-eq-Nil2 hd-append2 hd-conv-nth in-set-conv-decomp
insert-subset linorder-not-less list.simps(15) non-triangle not-less-eq not-less-iff-gr-or-eq
numeral-3-eq-3 same-set set-append snoc-eq-iff-butlast vts-is)
then have main-subset: set ?vtsp2  $\subseteq$  set (butlast vts)
  using subset3 subset4 by simp

have subset-p1: set ?vtsp1  $\subset$  set (butlast vts)
  using ij-prop distinct-butlast-vts
proof -

```

```

have card (set ?vtsp2) ≥ 3
  using polygon-p2 polygon-at-least-3-vertices by blast
moreover have set ?vtsp1 ∩ set ?vtsp2 = {vts!i, vts!j}
proof-
  have set ?vts2 ∩ set ?vts3 = {}
  by (metis append-take-drop-id diff-le-self distinct-append distinct-butlast-vts
set-take-disj-set-drop-if-distinct)
  moreover have set ?vts2 ∩ set ?vts1 = {}
  proof-
    have set ?vts2 ⊆ set (drop (i + 1) vts)
      by (metis add.commute drop-butlast in-set-butlastD in-set-takeD
plus-1-eq-Suc subset-code(1))
    moreover have set (drop (i + 1) vts) ∩ set ?vts1 ⊆ {last vts}
    proof-
      have set (drop (i + 1) (butlast vts)) ∩ set ?vts1 = {}
      by (simp add: Int-commute set-take-disj-set-drop-if-distinct dis-
tinct-butlast-vts)
    moreover have set (drop (i + 1) vts) = set (drop (i + 1) (butlast
vts)) ∪ {last vts}
    proof-
      have drop (i + 1) vts = (drop (i + 1) ((butlast vts) @ [last vts]))
      by (metis last-snoc vts-is)
      thus ?thesis using ij-prop by force
    qed
  ultimately show ?thesis by blast
  qed
  moreover have last vts ∉ set ?vts2
  by (metis card-set-len-butlast card-vts distinct-butlast-vts dual-order.strict-trans1
in-set-takeD index-nth-id last-snoc nth-butlast numeral-3-eq-3 set-drop-if-index vts-is
zero-less-Suc)
  ultimately show ?thesis by force
  qed
  moreover have vts!i ∈ set ?vtsp1 by (metis ij-prop list.set-intros(1)
nth-butlast)
  moreover have vts!j ∈ set ?vtsp1 using ij-prop nth-butlast by fastforce
  moreover have vts!i ∈ set ?vtsp2
  by (metis UnCI ij-prop list.set-intros(1) nth-butlast set-append)
  moreover have vts!j ∈ set ?vtsp2 using ij-prop nth-butlast by force
  moreover have set ?vtsp1 = set ?vts2 ∪ {vts!i, vts!j}
  by (smt (verit, ccfv-SIG) Un-insert-right empty-set ij-prop insert-absorb2
insert-commute list.simps(15) nth-butlast set-append)
  moreover have set ?vtsp2 = set ?vts1 ∪ set ?vts3 ∪ {vts!i, vts!j, vts!0}
  proof-
    have vts!i = (butlast vts)!i by (metis ij-prop nth-butlast)
    moreover have vts!j = (butlast vts)!j by (metis ij-prop nth-butlast)
    moreover have vts!0 = (butlast vts)!0
    by (metis ij-prop leD length-greater-0-conv nth-butlast take-all-iff
take-eq-Nil)
  ultimately show ?thesis by force

```

qed
moreover have $vts!0 \notin \text{set } ?vts2$
by (*metis distinct-butlast-vts in-set-conv-decomp in-set-takeD index-nth-id length-pos-if-in-set nth-butlast same-set set-drop-if-index vts-is zero-less-Suc*)
ultimately show *?thesis* **by** *blast*
qed
ultimately have $\text{card } (\text{set } ?vtsp2) > \text{card } (\text{set } ?vtsp1 \cap \text{set } ?vtsp2)$
by (*smt (verit, del-insts) card-length empty-set leI le-trans length-Cons list.simps(15) list.size(3) not-less-eq-eq numeral-3-eq-3*)
then have $\exists v. v \in \text{set } ?vtsp2 \wedge v \notin (\text{set } ?vtsp1 \cap \text{set } ?vtsp2)$
by (*smt (verit) Int-lower2 Orderings.order-eq-iff less-not-refl subset-code(1)*)
then obtain v **where** $v \in \text{set } ?vtsp2 - \text{set } ?vtsp1$ **by** *blast*
thus *?thesis*
by (*metis main-subset Diff-eq-empty-iff length-pos-if-in-set less-numeral-extra(3) list.set(1) list.size(3) psubsetI vtsp1-subset*)
qed
then have $\text{card } (\text{set } ?vtsp1) < \text{card } (\text{set } (\text{butlast } vts))$
using *card-subset-eq[OF finite-butlast]*
by (*meson finite-butlast psubset-card-mono*)
then have *card-lt-p1*: $\text{card } (\text{set } ?vtsp1) < \text{card } (\text{set } vts)$
using *same-set* **by** *argo*
have $\text{set } ?vtsp1 \subseteq \text{set } vts$
using *ij-prop*
using *same-set subset-p1* **by** *blast*
then have *all-integral-p1*: *all-integral* $?vtsp1$
using *less(4) unfolding all-integral-def*
by *blast*

obtain $p1' vtsp1'$ **where** *p1-rot*: *polygon-of* $p1' vtsp1'$
 $\wedge vtsp1'!0 \in \text{frontier } (\text{convex hull } (\text{set } vtsp1'))$
 $\wedge \text{path-image } p1' = \text{path-image } ?p1$
 $\wedge \text{all-integral } vtsp1'$
 $\wedge \text{set } vtsp1' = \text{set } ?vtsp1$
using *pick-rotate less polygon-p1 unfolding polygon-of-def*
using *all-integral-p1*
by *blast*

let $?I1' = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p1'\}$
let $?B1' = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p1'\}$

have *measure lebesgue* $(\text{path-inside } p1') = \text{real } ?I1' + \text{real } ?B1' / 2 - 1$
using *less(1) polygon-split card-lt-p1 p1-rot unfolding polygon-of-def* **by**
force
then have *indh1*: *Sigma-Algebra.measure lebesgue* $(\text{path-inside } ?p1) = \text{real } ?I1 + \text{real } ?B1 / 2 - 1$
using *p1-rot unfolding path-inside-def* **by** *metis*

have $vts ! (i+1) \notin \text{set } (\text{take } i (\text{butlast } vts))$
using *distinct-butlast-vts j-neq ij-prop*

proof–
have $i + 1 < \text{length } vts - 2$ **using** *distinct-butlast-vts j-neq ij-prop* **by** *fastforce*
then have $vts ! (i+1) = (\text{butlast } vts) ! (i+1)$ **by** (*simp add: nth-butlast*)
moreover then have $\forall j < i + 1. (\text{butlast } vts) ! j \neq (\text{butlast } vts) ! (i+1)$
using *distinct-butlast-vts distinct-nth-eq-iff ij-prop* **by** *fastforce*
moreover have $\text{set } (\text{take } i (\text{butlast } vts)) = \{vts!j \mid j. j < i\}$
proof–
have $\text{set } (\text{take } i (\text{butlast } vts)) \subseteq \{vts!j \mid j. j < i\}$
by (*smt (verit, ccfv-SIG) dual-order.strict-trans ij-prop in-set-conv-nth length-take mem-Collect-eq min.absorb4 nth-butlast nth-take subsetI*)
moreover have $\{vts!j \mid j. j < i\} \subseteq \text{set } (\text{take } i (\text{butlast } vts))$
by (*smt (verit, del-insts) dual-order.strict-trans ij-prop in-set-conv-nth length-take mem-Collect-eq min.absorb4 nth-butlast nth-take subsetI*)
ultimately show *?thesis* **by** *blast*
qed
ultimately show *?thesis*
by (*metis (no-types, lifting) add commute ij-prop in-set-conv-nth length-take min.absorb4 nth-take trans-less-add2*)
qed
moreover have $vts ! (i+1) \neq \text{butlast } vts ! i$
by (*metis (no-types, lifting) ij-prop add commute add-cancel-right-right distinct-butlast-vts distinct-nth-eq-iff less-trans-Suc nth-append plus-1-eq-Suc vts-is zero-neq-one*)
moreover have $vts ! (i+1) \neq \text{butlast } vts ! j$
by (*metis (no-types, lifting) add commute distinct-butlast-vts distinct-nth-eq-iff ij-prop j-neq less-trans-Suc nth-append plus-1-eq-Suc vts-is*)
ultimately have $vts ! (i+1) \notin \text{set } (\text{take } i (\text{butlast } vts))$ @
 $[\text{butlast } vts ! i, \text{butlast } vts ! j]$ **by** *force*
moreover have $vts ! (i+1) \notin \text{set } (\text{drop } (j - i) (\text{drop } (\text{Suc } i) (\text{butlast } vts)))$ @
 $[\text{butlast } vts ! 0]$
proof–
have $vts ! (i+1) \notin \text{set } (\text{drop } (j - i + \text{Suc } i) (\text{butlast } vts))$
by (*metis (no-types, lifting) add commute distinct-butlast-vts ij-prop index-nth-id less-add-same-cancel2 less-trans-Suc nth-append plus-1-eq-Suc set-drop-if-index vts-is zero-less-diff*)
moreover have $vts ! (i+1) \neq \text{butlast } vts ! 0$
by (*metis (no-types, lifting) ij-prop Nil-is-append-conv add commute distinct-butlast-vts distinct-nth-eq-iff length-greater-0-conv less-trans-Suc list.discI nat.distinct(1) nth-append plus-1-eq-Suc same-set set-empty vts-is*)
ultimately show *?thesis* **by** *simp*
qed
ultimately have $vts ! (i+1) \notin \text{set } (\text{take } i (\text{butlast } vts))$ @
 $[\text{butlast } vts ! i, \text{butlast } vts ! j]$ @
 $\text{drop } (j - i) (\text{drop } (\text{Suc } i) (\text{butlast } vts))$ @ $[\text{butlast } vts ! 0]$
by *auto*
then have $\text{subset-butlast-p2: set } ?vtsp2 \subset \text{set } (\text{butlast } vts)$
using *main-subset ij-prop*
by (*metis (no-types, lifting) antisym-conv2 length-butlast less-diff-conv*)

```

nth-mem same-set)
  then have card-lt-p2: card (set ?vtsp2) < card (set vts)
    using card-subset-eq[OF finite-butlast]
    by (metis finite-butlast psubset-card-mono same-set)
  have subset-p2: set ?vtsp2  $\subset$  set vts
    using subset-butlast-p2 same-set
    by presburger
  then have all-integral-p2: all-integral ?vtsp2
    using less(4) unfolding all-integral-def
    by blast

  let ?p2 = make-polygonal-path (take i (butlast vts) @ [butlast vts ! i, butlast
vts ! j] @
    drop (j - i) (drop (Suc i) (butlast vts)) @ [butlast vts ! 0])
  let ?I2 = card {x. integral-vec x  $\wedge$  x  $\in$  path-inside ?p2}
  let ?B2 = card {x. integral-vec x  $\wedge$  x  $\in$  path-image ?p2}
  have polygon-p2: polygon ?p2
    using polygon-split unfolding is-polygon-split-def by metis

  have vtsp2-0: ?vtsp2!0  $\in$  frontier (convex hull (set ?vtsp2))
  proof-
    have ?vtsp2!0 = vts!0
      by (metis (no-types, lifting) append-Cons ij-prop length-greater-0-conv
less-nat-zero-code nat-neq-iff nth-append nth-append-length nth-butlast nth-take take-eq-Nil)
    then have ?vtsp2!0  $\in$  frontier (convex hull (set vts)) using less by argo
    moreover have ?vtsp2!0  $\in$  (convex hull (set ?vtsp2))
      by (meson append-is-Nil-conv hull-inc length-greater-0-conv neq-Nil-conv
nth-mem)
    moreover have convex hull (set ?vtsp2)  $\subseteq$  convex hull (set vts)
      by (metis hull-mono main-subset same-set)
    ultimately show ?thesis using in-frontier-in-subset by blast
  qed

  have indh2: Sigma-Algebra.measure lebesgue (path-inside ?p2) = real ?I2 +
real ?B2 / 2 - 1
    using less(1)[OF card-lt-p2 polygon-p2 - all-integral-p2 - - vtsp2-0] poly-
gon-split
    by blast

  have all-integral (butlast vts)  $\implies$ 
    Sigma-Algebra.measure lebesgue (path-inside p) = real (card {x. integral-vec
x  $\wedge$  x  $\in$  path-inside p}) + real (card {x. integral-vec x  $\wedge$  x  $\in$  path-image p}) / 2
- 1
    using pick-split-union
      [OF polygon-split, of ?vts1 ?vts2 ?vts3 butlast vts ! i butlast vts ! j p ?p1
?p2 ?I1 ?B1 ?I2 ?B2]
    using indh1 indh2 p-is
    by blast
  then have ?case

```

```

    using less(4-6) unfolding all-integral-def
    using same-set by presburger
  } moreover
  { assume non-convex:  $\neg$  (convex (path-image p  $\cup$  path-inside p))
    let ?vts-ch = set vts  $\cap$  frontier (convex hull (set vts))
    have finite-vts: finite (set vts)
      using less
      by force
    have subset-ch: ?vts-ch  $\subset$  set vts
      using vts-subset-frontier
      using less.premis(1) less.premis(2) non-convex polygon-of-def by blast
    then have card-ch: card (?vts-ch) < card (set vts)
      using finite-vts
      by (simp add: psubset-card-mono)

    let ?vts-ch-list = filter ( $\lambda v. v \in ?vts-ch$ ) vts

    let ?r-idx = min-index-not-in-set vts ?vts-ch
    let ?r = ?r-idx - 1
    let ?rotated-vts = rotate-polygon-vertices vts ?r
    let ?pr = make-polygonal-path ?rotated-vts

    have subset-ch-list: set ?vts-ch-list  $\subset$  set vts using subset-ch by auto
    then have r-defined: index-not-in-set vts ?vts-ch ?r-idx
       $\wedge$  ( $\forall j < ?r-idx. \neg$  index-not-in-set vts ?vts-ch j)
      using min-index-not-in-set-defined[of ?vts-ch vts] by fastforce

    have pr-image: path-image p = path-image ?pr
      using polygon-vts-arb-rotation less by blast
    then have measure lebesgue (path-inside ?pr) = measure lebesgue (path-inside
  p)
      unfolding path-inside-def by presburger
    have rotated-vts-set: set ?rotated-vts = set vts
      using less.premis(1) less.premis(2) rotate-polygon-vertices-same-set by auto
    then have card (set ?rotated-vts) = card (set vts) by argo
    have polygon-rotation: polygon ?pr using rotation-is-polygon less by blast

    let ?pocket-path-vts = construct-pocket-0 ?rotated-vts ?vts-ch

    let ?a = hd ?pocket-path-vts
    let ?b = last ?pocket-path-vts
    let ?l = linepath ?a ?b

    have vts!0  $\in$  ?vts-ch
      by (metis IntI length-greater-0-conv less.premis(6) nth-mem snoc-eq-iff-butlast
    vts-is)
    then have vts-r: vts! ?r  $\in$  ?vts-ch

```

using *min-index-not-in-set-0 subset-ch* **by** *presburger*
moreover have *rotated-0: ?rotated-vts!0 = vts! ?r*
using *rotated-polygon-vertices*[of *?rotated-vts vts ?r ?r*]
by (*metis (no-types, lifting) Suc-1 Suc-leI card-gt-0-iff card-set-len-butlast*
diff-is-0-eq' finite-vts hd-conv-nth index-not-in-set-def le-refl length-butlast less-imp-diff-less
mem-Collect-eq r-defined set-empty snoc-eq-iff-butlast vts-is zero-less-diff)
ultimately have *rotated-0-in: ?rotated-vts!0 ∈ ?vts-ch* **by** *presburger*
then have *b-in: ?b ∈ set vts*
using *construct-pocket-0-last-in-set*[of *?rotated-vts ?vts-ch*]
by (*smt (verit, ccfv-threshold) Int-iff One-nat-def closed-path-def Suc-leI*
card-0-eq card-set-len-butlast empty-iff finite-vts last-conv-nth last-in-set last-tl length-butlast
length-greater-0-conv length-tl list.size(3) polygon-def polygon-pathfinish polygon-pathstart
polygon-rotation rotate-polygon-vertices-same-length set-empty)

have $2 \leq \text{card } ?vts\text{-ch}$
using *convex-hull-two-vts-on-frontier*
by (*metis One-nat-def Suc-1 add-leD2 card-vts numeral-3-eq-3 plus-1-eq-Suc*)
moreover have $?vts\text{-ch} \subseteq \text{set } ?rotated\text{-vts}$
using *less.premis(1) less.premis(2) rotate-polygon-vertices-same-set* **by** *force*
moreover have *distinct (butlast ?rotated-vts)*
using *polygon-def polygon-rotation simple-polygonal-path-vts-distinct* **by** *blast*
moreover have *hd-last-rotated: hd ?rotated-vts = last ?rotated-vts*
by (*metis have-wraparound-vertex hd-conv-nth polygon-rotation snoc-eq-iff-butlast*)
ultimately have $a \neq b: ?a \neq ?b$
using *construct-pocket-0-first-last-distinct*
by (*smt (verit) Collect-cong Int-def mem-Collect-eq set-filter*)

let $?pocket\text{-vts} = ?pocket\text{-path-vts} @ [?rotated\text{-vts!}0]$

let $?pocket\text{-good-path-vts} = \text{tl (butlast } ?pocket\text{-path-vts)}$

let $?filled\text{-vts} = \text{fill-pocket-0 } ?rotated\text{-vts (length } ?pocket\text{-path-vts)}$

let $?filled\text{-vts-tl} = \text{tl } ?filled\text{-vts}$

let $?filled\text{-p-tl} = \text{make-polygonal-path } ?filled\text{-vts-tl}$

let $?filled\text{-p} = \text{make-polygonal-path } ?filled\text{-vts}$

let $?pocket\text{-path} = \text{make-polygonal-path } ?pocket\text{-path-vts}$

let $?pocket = \text{make-polygonal-path } ?pocket\text{-vts}$

have *non-convex-rot: $\neg \text{convex (path-image } ?pr \cup \text{path-inside } ?pr)$*

using *non-convex* **by** (*simp add: path-inside-def pr-image*)

have $0: ?rotated\text{-vts!}0 \in \text{frontier (convex hull (set } ?rotated\text{-vts))}$

using *less.premis(1) less.premis(2) rotate-polygon-vertices-same-set* *rotated-0-in* **by** *fastforce*

have $1: ?rotated\text{-vts!}1 \notin \text{frontier (convex hull (set } ?rotated\text{-vts))}$

proof—

```

have ?rotated-vts!1 = vts!(?r + 1)
  using rotated-polygon-vertices[of ?rotated-vts vts ?r ?r + 1]
  by (smt (verit, ccfv-threshold) Suc-1 Suc-leI card-gt-0-iff card-set-len-butlast
diff-is-0-eq' finite-vts hd-conv-nth index-not-in-set-def le-refl length-butlast less-imp-diff-less
mem-Collect-eq r-defined set-empty snoc-eq-iff-butlast vts-is zero-less-diff Suc-diff-Suc
add commute add-diff-cancel-left' bot-nat-0.not-eq-extremum less-imp-le-nat plus-1-eq-Suc)
  also have ...  $\notin$  frontier (convex hull (set ?rotated-vts))
    using r-defined unfolding index-not-in-set-def
    by (smt (verit, best) Int-iff Suc-leI add commute add-diff-inverse-nat
bot-nat-0.not-eq-extremum diff-is-0-eq' mem-Collect-eq nat-less-le nth-mem plus-1-eq-Suc
rotated-vts-set vts-r zero-less-diff)
  finally show ?thesis .
qed
then have split:
  is-polygon-split-path (butlast ?filled-vts) 0 1 ?pocket-good-path-vts
  and polygon-filled-p: polygon ?filled-p
  and polygon-pocket: polygon ?pocket
  and pocket-path-vts-card: card (set ?pocket-path-vts) < card (set vts)
  and filled-vts-card: card (set ?filled-vts) < card (set vts)
    using pocket-path-good[OF - 0 1 non-convex-rot] polygon-rotation ro-
tated-vts-set apply argo
    using pocket-path-good[OF - 0 1 non-convex-rot] polygon-rotation ro-
tated-vts-set apply argo
    using pocket-path-good[OF - 0 1 non-convex-rot] polygon-rotation ro-
tated-vts-set
  apply (metis add-gr-0 construct-pocket-0-def nth-take zero-less-one)
    using pocket-path-good[OF - 0 1 non-convex-rot] polygon-rotation ro-
tated-vts-set apply argo
    using pocket-path-good[OF - 0 1 non-convex-rot] polygon-rotation ro-
tated-vts-set by argo

  have vts-0-frontier: ?rotated-vts!0  $\in$  frontier (convex hull (set vts))
    using rotated-0-in by simp
  have filled-0: ?filled-vts!0 = ?rotated-vts!0
  by (metis convex-hull-empty empty-set fill-pocket-0-def frontier-empty hd-conv-nth
length-pos-if-in-set less.premis(6) less-numeral-extra(3) list.size(3) nth-Cons-0 ro-
tated-vts-set)
  have pocket-0: ?pocket-vts!0 = ?rotated-vts!0
    unfolding construct-pocket-0-def
    by (simp add: less-numeral-extra(1) nth-append trans-less-add2)

  have subset-pocket-path-vts: set ?pocket-path-vts  $\subseteq$  set vts
    using construct-pocket-0-subset-vts
  by (metis construct-pocket-0-def less.premis(1) less.premis(2) rotate-polygon-vertices-same-set
set-take-subset)
  moreover have set ?pocket-good-path-vts  $\subseteq$  set ?pocket-path-vts
  by (smt (verit, best) butlast-conv-take list.exhaust-sel list.sel(2) set-subset-Cons
set-take-subset subset-trans)
  ultimately have subset-pocket-good-path: set ?pocket-good-path-vts  $\subseteq$  set vts

```

by *blast*
then have *subset-pocket*: $set\ ?pocket\ vts \subseteq set\ vts$
by (*metis* (*mono-tags*, *lifting*) *have-wraparound-vertex* *less.premis*(1) *less.premis*(2)
polygon-rotation *rotate-polygon-vertices-same-set* *set-append* *subset-code*(1) *subset-pocket-path-vts*
sup.bounded-iff)
have $set\ ?filled\ vts \subseteq set\ ?rotated\ vts$
unfolding *fill-pocket-0-def*
by (*metis* *b-in* *hd-in-set* *insert-subset* *length-pos-if-in-set* *less-numeral-extra*(3)
list.simps(15) *list.size*(3) *rotated-vts-set* *set-drop-subset*)
then have *subset-filled*: $set\ ?filled\ vts \subseteq set\ vts$
using *rotated-vts-set* **by** *blast*

have *taut1*: $?filled\ p = make\ polygonal\ path\ ?filled\ vts$ **by** *blast*
have *all-integral-filled-vts*: $all\ integral\ ?filled\ vts$
using *subset-filled* **less** **by** (*meson* *all-integral-def* *subset-iff*)
have *taut2*: $card\ (integral\ inside\ ?filled\ p) = card\ \{x.\ integral\ vec\ x \wedge x \in$
*path-inside\ ?filled\ p\}
unfolding *integral-inside* **by** *blast*
have *taut3*: $card\ (integral\ boundary\ ?filled\ p) = card\ \{x.\ integral\ vec\ x \wedge x \in$
*path-image\ ?filled\ p\}
unfolding *integral-boundary* **by** *blast*
have *filled-vts-0-frontier*: $?filled\ vts!0 \in frontier\ (convex\ hull\ (set\ ?filled\ vts))$
proof–
have $?filled\ vts!0 \in frontier\ (convex\ hull\ set\ vts)$
using *filled-0* *vts-0-frontier* **by** *presburger*
moreover have $?filled\ vts!0 \in convex\ hull\ (set\ ?filled\ vts)$
by (*metis* *have-wraparound-vertex* *hull-inc* *in-set-conv-decomp* *polygon-filled-p*)
moreover have $set\ ?filled\ vts \subseteq set\ vts$ **using** *subset-filled* **by** *force*
ultimately show *?thesis* **using** *in-frontier-in-subset-convex-hull* **by** *blast*
qed

have *ih-filled*: $measure\ lebesgue\ (path\ inside\ ?filled\ p)$
 $= card\ (integral\ inside\ ?filled\ p) + ((card\ (integral\ boundary\ ?filled\ p)) /$
 $2) - 1$
using *less*(1)[*OF* *filled-vts-card* *polygon-filled-p* *taut1* *all-integral-filled-vts*
taut2 *taut3* *filled-vts-0-frontier*]
by *blast*

have $set\ ?pocket\ path\ vts \subset set\ vts$
using *pocket-path-vts-card* *subset-pocket-path-vts* **by** *force*
moreover have *pocket-path-set*: $set\ ?pocket\ path\ vts = set\ ?pocket\ vts$
by (*smt* (*verit*) *Nil-is-append-conv* *rotated-0* *a-neq-b* *append-Cons* *append-Nil*
hd-Nil-eq-last *hd-append2* *hd-conv-nth* *hd-in-set* *insert-absorb* *list.simps*(15) *pocket-0*
rev-append *set-append* *set-rev*)
ultimately have $set\ ?pocket\ vts \subset set\ vts$ **by** *blast*
then have *pocket-vts-card*: $card\ (set\ ?pocket\ vts) < card\ (set\ vts)$
by (*meson* *finite-vts* *psubset-card-mono*)**

```

have all-integral-pocket-vts: all-integral ?pocket-vts
  using subset-pocket less unfolding all-integral-def by blast
have taut1: ?pocket = make-polygonal-path ?pocket-vts by blast
have taut2: card (integral-inside ?pocket) = card {x. integral-vec x ∧ x ∈
path-inside ?pocket}
  unfolding integral-inside by blast
have taut3: card (integral-boundary ?pocket) = card {x. integral-vec x ∧ x ∈
path-image ?pocket}
  unfolding integral-boundary by blast
have pocket-vts-0-frontier: ?pocket-vts!0 ∈ frontier (convex hull (set ?pocket-vts))
proof–
  have ?pocket-vts!0 ∈ frontier (convex hull set vts)
    using pocket-0 vts-0-frontier by presburger
  moreover have ?pocket-vts!0 ∈ convex hull (set ?pocket-vts)
    by (smt (verit, del-insts) hull-inc in-set-conv-decomp pocket-0)
  moreover have set ?pocket-vts ⊆ set vts using subset-pocket by force
  ultimately show ?thesis using in-frontier-in-subset-convex-hull by blast
qed

have ih-pocket: measure lebesgue (path-inside ?pocket) = card (integral-inside
?pocket) + ((card (integral-boundary ?pocket)) / 2) – 1
  using less(1)[OF pocket-vts-card polygon-pocket taut1 all-integral-pocket-vts
taut2 taut3 pocket-vts-0-frontier]
  by blast

```

```

let ?i = 0::nat
let ?j = 1::nat
let ?vts = butlast ?filled-vts
let ?vts1 = []
let ?vts2 = []
let ?vts3 = butlast (drop 2 ?filled-vts)
let ?cutvts = ?pocket-good-path-vts
let ?p = ?filled-p
let ?p1 = make-polygonal-path (?a # ?vts2 @ [?b] @ rev ?cutvts @ [?a])
let ?p2 = ?pr
let ?I1 = card {x. integral-vec x ∧ x ∈ path-inside ?p1}
let ?B1 = card {x. integral-vec x ∧ x ∈ path-image ?p1}
let ?I2 = card {x. integral-vec x ∧ x ∈ path-inside ?p2}
let ?B2 = card {x. integral-vec x ∧ x ∈ path-image ?p2}
let ?I = card {x. integral-vec x ∧ x ∈ path-inside ?p}
let ?B = card {x. integral-vec x ∧ x ∈ path-image ?p}

```

```

have rev ?pocket-vts = (?a # ?vts2 @ [?b] @ rev ?cutvts @ [?a])
  by (smt (verit) a-neq-b append-Nil append-butlast-last-id hd-Nil-eq-last
hd-append2 hd-conv-nth last-conv-nth length-butlast list.collapse list.size(3) pocket-0
rev.simps(2) rev-append rev-rev-ident snoc-eq-iff-butlast)
then have pocket-rev-image: path-image ?pocket = path-image ?p1
  using polygon-at-least-3-vertices polygon-pocket card-length
  by (smt (verit, best) One-nat-def Suc-1 le-add2 le-trans numeral-3-eq-3)

```

```

plus-1-eq-Suc rev-pts-path-image polygon-at-least-3-vertices polygon-pocket card-length)
  then have pocket-rev-inside: path-inside ?pocket = path-inside ?p1
    unfolding path-inside-def by argo

  have split': is-polygon-split-path ?pts ?i ?j ?cutpts using split by blast
  have 0: ?pts1 = take ?i ?pts by auto
  have 1: ?pts2 = take (?j - ?i - 1) (drop (Suc ?i) ?pts) by simp
  have 2: ?pts3 = drop (?j - ?i) (drop (Suc ?i) ?pts)
  by (metis (no-types, lifting) One-nat-def Suc-1 diff-zero drop-butlast drop-drop
plus-1-eq-Suc)
  have 3: ?a = ?pts ! ?i
  by (smt (z3) Nil-is-append-conv pocket-path-set filled-0 hd-conv-nth is-polygon-split-path-def
length-greater-0-conv list.distinct(1) nth-append nth-butlast pocket-0 set-empty split')
  have 4: ?b = ?pts ! ?j
  proof-
    have ?b = ?filled-pts!1
      unfolding construct-pocket-0-def fill-pocket-0-def
      by (smt (z3) Suc-eq-plus1 a-neq-b construct-pocket-0-def diff-Suc-1
diff-is-0-eq' drop-eq-Nil hd-conv-nth hd-drop-conv-nth hd-last-rotated last-conv-nth
length-take linorder-not-less min.absorb4 nat-le-linear not-less-eq-eq nth-Cons' nth-take
one-neq-zero take-all-iff take-eq-Nil)
    thus ?thesis by (metis is-polygon-split-path-def nth-butlast split')
  qed
  have 5: ?pocket-path = make-polygonal-path (?a # ?cutpts @ [?b])
    by (smt (verit, ccfv-SIG) a-neq-b butlast.simps(2) butlast-tl hd-Cons-tl
hd-Nil-eq-last last.simps snoc-eq-iff-butlast)
  have 6: ?p = make-polygonal-path (?pts @ [?pts!0])
  by (metis (no-types, lifting) butlast-conv-take have-wraparound-vertex is-polygon-split-path-def
nth-butlast polygon-filled-p split')
  have 7: ?p1 = make-polygonal-path (?a # ?pts2 @ [?b] @ rev ?cutpts @ [?a])
  by blast
  have 8: ?p2 = make-polygonal-path (?pts1 @ ([?a] @ ?cutpts @ [?b]) @ ?pts3
@ [?pts!0])
  proof-
    have ?rotated-pts = ?pts1 @ ([?a] @ ?cutpts @ [?b]) @ ?pts3 @ [?pts!0]
      unfolding construct-pocket-0-def fill-pocket-0-def
    by (smt (verit) 3 Suc-1 hd-last-rotated a-neq-b append-Cons append-Nil ap-
pend-butlast-last-id append-take-drop-id construct-pocket-0-def drop-Suc drop-drop
drop-eq-Nil fill-pocket-0-def hd-Nil-eq-last hd-append2 hd-conv-nth last-conv-nth last-drop
length-Cons length-take length-tl linorder-not-less list.collapse list.sel(3) list.size(3)
min.absorb4 plus-1-eq-Suc take-all-iff)
    thus ?thesis by argo
  qed
  have 9: ?I1 = card {x. integral-vec x ∧ x ∈ path-inside ?p1} by blast
  have 10: ?B1 = card {x. integral-vec x ∧ x ∈ path-image ?p1} by blast
  have 11: ?I2 = card {x. integral-vec x ∧ x ∈ path-inside ?p2} by blast
  have 12: ?B2 = card {x. integral-vec x ∧ x ∈ path-image ?p2} by blast
  have 13: ?I = card {x. integral-vec x ∧ x ∈ path-inside ?p} by blast
  have 14: ?B = card {x. integral-vec x ∧ x ∈ path-image ?p} by blast

```

```

have 15: all-integral ?vts
  using subset-filled less
  unfolding all-integral-def
by (metis (no-types, lifting) all-integral-def all-integral-filled-vts in-set-butlastD)
have 16: measure lebesgue (path-inside ?p) = ?I + ?B/2 - 1
  using ih-filled unfolding integral-inside integral-boundary by blast
have 17: measure lebesgue (path-inside ?p1) = ?I1 + ?B1/2 - 1
using ih-pocket unfolding integral-inside integral-boundary using pocket-rev-image
pocket-rev-inside by force
have measure lebesgue (path-inside ?p2) = ?I2 + ?B2/2 - 1
  using pick-split-path-union-main(3)
  [OF split' 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17] less(5-6) by blast
moreover have  $?I2 = I$  using less(5) pr-image path-inside-def by presburger
moreover have  $?B2 = B$  using less(6) pr-image path-image-def by pres-
burger
  ultimately have ?case by (simp add: path-inside-def pocket-rev-inside
pr-image)
}
ultimately have ?case by blast
}
ultimately show ?case using card-vts by linarith
qed

```

theorem *pick*:

```

fixes  $p :: R\text{-to-}R^2$ 
assumes polygon p
assumes  $p = \text{make-polygonal-path } vts$ 
assumes all-integral vts
assumes  $I = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-inside } p\}$ 
assumes  $B = \text{card } \{x. \text{integral-vec } x \wedge x \in \text{path-image } p\}$ 
shows measure lebesgue (path-inside p) = I + B/2 - 1
proof -
obtain  $p' vts'$  where polygon-of p' vts'
   $\wedge vts' \neq \emptyset \in \text{frontier } (\text{convex hull } (\text{set } vts'))$ 
   $\wedge \text{path-image } p' = \text{path-image } p$ 
   $\wedge \text{all-integral } vts'$ 
   $\wedge \text{set } vts' = \text{set } vts$ 
using pick-rotate assms unfolding polygon-of-def by blast
thus ?thesis using assms pick-unrotated unfolding path-inside-def polygon-of-def
by fastforce
qed

```

end

References

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- [2] J. Harrison. A formal proof of Pick's theorem. *Math. Struct. Comput. Sci.*, 21(4):715–729, 2011.