

Pell's Equation

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Abstract

This article gives the basic theory of Pell's equation $x^2 = 1 + Dy^2$, where $D \in \mathbb{N}$ is a parameter and x, y are integer variables.

The main result that is proven is the following: If D is not a perfect square, then there exists a *fundamental solution* (x_0, y_0) that is not the trivial solution $(1, 0)$ and which generates all other solutions (x, y) in the sense that there exists some $n \in \mathbb{N}$ such that $|x| + |y|\sqrt{D} = (x_0 + y_0\sqrt{D})^n$. This also implies that the set of solutions is infinite, and it gives us an explicit and executable characterisation of all the solutions.

Based on this, simple executable algorithms for computing the fundamental solution and the infinite sequence of all non-negative solutions are also provided.

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```

theory Efficient-Discrete-Sqrt
imports
  Complex-Main
  HOL-Computational-Algebra.Computational-Algebra
  HOL-Library.Discrete
  HOL-Library.Tree
  HOL-Library.IArray
begin

```

1 Efficient Algorithms for the Square Root on \mathbb{N}

1.1 A Discrete Variant of Heron's Algorithm

An algorithm for calculating the discrete square root, taken from Cohen [2]. This algorithm is essentially a discretised variant of Heron's method or Newton's method specialised to the square root function.

```

lemma sqrt-eq-floor-sqrt: Discrete.sqrt n = nat [sqrt n]
proof -
  have real ((nat [sqrt n])2) = (real (nat [sqrt n]))2
    by simp
  also have ... ≤ sqrt (real n) ^ 2
    by (intro power-mono) auto
  also have ... = real n by simp
  finally have (nat [sqrt n])2 ≤ n
    by (simp only: of-nat-le-iff)
  moreover have n < (Suc (nat [sqrt n]))2 proof -
    have (1 + [sqrt n])2 > n
      using floor-correct[of sqrt n] real-le-rsqrt[of 1 + [sqrt n] n]
        of-int-less-iff[of n (1 + [sqrt n])2] not-le
      by fastforce
    then show ?thesis
      using le-nat-floor[of Suc (nat [sqrt n]) sqrt n]
        of-nat-le-iff[of (Suc (nat [sqrt n]))2 n] real-le-rsqrt[of - n] not-le
      by fastforce
    qed
  ultimately show ?thesis using sqrt-unique by fast
qed

```

```

fun newton-sqrt-aux :: nat ⇒ nat ⇒ nat where
  newton-sqrt-aux x n =
    (let y = (x + n div x) div 2
     in if y < x then newton-sqrt-aux y n else x)

```

```

declare newton-sqrt-aux.simps [simp del]

```

```

lemma newton-sqrt-aux-simps:
  (x + n div x) div 2 < x ⇒ newton-sqrt-aux x n = newton-sqrt-aux ((x + n div

```

$x \text{ div } 2 \text{) } n$
 $(x + n \text{ div } x) \text{ div } 2 \geq x \implies \text{newton-sqrt-aux } x \ n = x$
by (subst newton-sqrt-aux.simps; simp add: Let-def)+

lemma heron-step-real: $\llbracket t > 0; n \geq 0 \rrbracket \implies (t + n/t) / 2 \geq \text{sqrt } n$
using arith-geo-mean-sqrt[of t n/t] **by** simp

lemma heron-step-div-eq-floored:

$(t::\text{nat}) > 0 \implies (t + (n::\text{nat}) \text{ div } t) \text{ div } 2 = \text{nat } \lfloor (t + n/t) / 2 \rfloor$

proof –

assume $t > 0$

then have $\lfloor (t + n/t) / 2 \rfloor = \lfloor (t*t + n) / (2*t) \rfloor$

by (simp add: mult-divide-mult-cancel-right[of t t + n/t 2, symmetric] algebra-simps)

also have $\dots = (t*t + n) \text{ div } (2*t)$

using floor-divide-of-nat-eq **by** blast

also have $\dots = (t*t + n) \text{ div } t \text{ div } 2$

by (simp add: Divides.div-mult2-eq mult commute)

also have $\dots = (t + n \text{ div } t) \text{ div } 2$

by (simp add: $\langle 0 < t \rangle$ power2-eq-square)

finally show ?thesis **by** simp

qed

lemma heron-step: $t > 0 \implies (t + n \text{ div } t) \text{ div } 2 \geq \text{Discrete.sqrt } n$

proof –

assume $t > 0$

have $\text{Discrete.sqrt } n = \text{nat } \lfloor \text{sqrt } n \rfloor$ **by** (rule sqrt-eq-floor-sqrt)

also have $\dots \leq \text{nat } \lfloor (t + n/t) / 2 \rfloor$

using heron-step-real[of t n] $\langle t > 0 \rangle$ **by** linarith

also have $\dots = (t + n \text{ div } t) \text{ div } 2$

using heron-step-div-eq-floored[OF $\langle t > 0 \rangle$] **by** simp

finally show ?thesis .

qed

lemma newton-sqrt-aux-correct:

assumes $x \geq \text{Discrete.sqrt } n$

shows $\text{newton-sqrt-aux } x \ n = \text{Discrete.sqrt } n$

using assms

proof (induction x n rule: newton-sqrt-aux.induct)

case (1 x n)

show ?case

proof (cases x = Discrete.sqrt n)

case True

then have $(x \wedge 2) \text{ div } x \leq n \text{ div } x$ **by** (intro div-le-mono) simp-all

also have $(x \wedge 2) \text{ div } x = x$ **by** (simp add: power2-eq-square)

finally have $(x + n \text{ div } x) \text{ div } 2 \geq x$ **by** linarith

with True **show** ?thesis **by** (auto simp: newton-sqrt-aux-simps)

next

case False

```

with 1.premis have x-gt-sqrt:  $x > \text{Discrete.sqrt } n$  by auto
with Discrete.le-sqrt-iff[of x n] have  $n < x^2$  by simp
have  $x * (n \text{ div } x) \leq n$  using mult-div-mod-eq[of x n] by linarith
also have  $\dots < x^2$  using Discrete.le-sqrt-iff[of x n] and x-gt-sqrt by simp
also have  $\dots = x * x$  by (simp add: power2-eq-square)
finally have  $n \text{ div } x < x$  by (subst (asm) mult-less-cancel1) auto
then have step-decreasing:  $(x + n \text{ div } x) \text{ div } 2 < x$  by linarith
with x-gt-sqrt have step-ge-sqrt:  $(x + n \text{ div } x) \text{ div } 2 \geq \text{Discrete.sqrt } n$ 
  by (simp add: heron-step)
from step-decreasing have newton-sqrt-aux x n = newton-sqrt-aux ((x + n div
x) div 2) n
  by (simp add: newton-sqrt-aux-simps)
also have  $\dots = \text{Discrete.sqrt } n$ 
  by (intro 1.IH step-decreasing step-ge-sqrt) simp-all
finally show ?thesis .
qed
qed

```

```

definition newton-sqrt :: nat  $\Rightarrow$  nat where
  newton-sqrt n = newton-sqrt-aux n n

```

```

declare Discrete.sqrt-code [code del]

```

```

theorem Discrete-sqrt-eq-newton-sqrt [code]:  $\text{Discrete.sqrt } n = \text{newton-sqrt } n$ 
  unfolding newton-sqrt-def by (simp add: newton-sqrt-aux-correct Discrete.sqrt-le)

```

1.2 Square Testing

Next, we implement an algorithm to determine whether a given natural number is a perfect square, as described by Cohen [2]. Essentially, the number first determines whether the number is a square. Essentially

```

definition q11 :: nat set
  where q11 = {0, 1, 3, 4, 5, 9}
definition q63 :: nat set
  where q63 = {0, 1, 4, 7, 9, 16, 28, 18, 22, 25, 36, 58, 46, 49, 37, 43}
definition q64 :: nat set
  where q64 = {0, 1, 4, 9, 16, 17, 25, 36, 33, 49, 41, 57}
definition q65 :: nat set
  where q65 = {0, 1, 4, 10, 14, 9, 16, 26, 30, 25, 29, 40, 56, 36, 49, 61, 35,
51, 39, 55, 64}

```

```

definition q11-array where
  q11-array = IArray [True, True, False, True, True, True, False, False, False, True, False]

```

```

definition q63-array where
  q63-array = IArray [True, True, False, False, True, False, False, True, False, True, False, False,
False, False, False, False, True, False, True, False, False, False, True, False, False, True, False,
False, True, False, False, False, False, False, False, False, True, True, False, False, False, False,

```


$(n \bmod 64 \in q64 \wedge (\text{let } r = n \bmod 45045 \text{ in } r \bmod 63 \in q63 \wedge r \bmod 65 \in q65 \wedge r \bmod 11 \in q11 \wedge n = (\text{Discrete.sqrt } n)^2))$

lemma *square-test-code* [code]:

square-test $n =$
 $(\text{IArray.sub } q64\text{-array } (n \bmod 64) \wedge (\text{let } r = n \bmod 45045 \text{ in } \text{IArray.sub } q63\text{-array } (r \bmod 63) \wedge \text{IArray.sub } q65\text{-array } (r \bmod 65) \wedge \text{IArray.sub } q11\text{-array } (r \bmod 11) \wedge n = (\text{Discrete.sqrt } n)^2))$
using *in-q11-code* [symmetric] *in-q63-code* [symmetric]
in-q64-code [symmetric] *in-q65-code* [symmetric]
by (*simp add: Let-def square-test-def*)

lemma *square-mod-lower*: $m > 0 \implies (q^2 :: \text{nat}) \bmod m = a \implies \exists q' < m. q'^2 \bmod m = a$

using *mod-less-divisor mod-mod-trivial power-mod* **by** *blast*

lemma *q11-upto-def*: $q11 = (\lambda k. k^2 \bmod 11) \text{ ‘ } \{..<11\}$
by (*simp add: q11-def lessThan-nat-numeral lessThan-Suc insert-commute*)

lemma *q11-infinite-def*: $q11 = (\lambda k. k^2 \bmod 11) \text{ ‘ } \{0..\}$
unfolding *q11-upto-def image-def* **proof** (*auto, goal-cases*)
case (*1 xa*)
show *?case*
using *square-mod-lower[of 11 xa xa^2 mod 11]*
ex-nat-less-eq[of 11 λx. xa^2 mod 11 = x^2 mod 11]
by *auto*

qed

lemma *q63-upto-def*: $q63 = (\lambda k. k^2 \bmod 63) \text{ ‘ } \{..<63\}$
by (*simp add: q63-def lessThan-nat-numeral lessThan-Suc insert-commute*)

lemma *q63-infinite-def*: $q63 = (\lambda k. k^2 \bmod 63) \text{ ‘ } \{0..\}$
unfolding *q63-upto-def image-def* **proof** (*auto, goal-cases*)
case (*1 xa*)
show *?case*
using *square-mod-lower[of 63 xa xa^2 mod 63]*
ex-nat-less-eq[of 63 λx. xa^2 mod 63 = x^2 mod 63]
by *auto*

qed

lemma *q64-upto-def*: $q64 = (\lambda k. k^2 \bmod 64) \text{ ‘ } \{..<64\}$
by (*simp add: q64-def lessThan-nat-numeral lessThan-Suc insert-commute*)

lemma *q64-infinite-def*: $q64 = (\lambda k. k^2 \bmod 64) \text{ ‘ } \{0..\}$
unfolding *q64-upto-def image-def* **proof** (*auto, goal-cases*)
case (*1 xa*)
show *?case*

```

using square-mod-lower[of 64 xa xa2 mod 64]
      ex-nat-less-eq[of 64 λx. xa2 mod 64 = x2 mod 64]
by auto
qed

lemma q65-upto-def: q65 = (λk. k2 mod 65) ‘{..65}
by (simp add: q65-def lessThan-nat-numeral lessThan-Suc insert-commute)

lemma q65-infinite-def: q65 = (λk. k2 mod 65) ‘{0..}
unfolding q65-upto-def image-def proof (auto, goal-cases)
case (1 xa)
show ?case
      using square-mod-lower[of 65 xa xa2 mod 65]
            ex-nat-less-eq[of 65 λx. xa2 mod 65 = x2 mod 65]
by auto
qed

lemma square-mod-existence:
  fixes n k :: nat
  assumes ∃ q. q2 = n
  shows ∃ q. n mod k = q2 mod k
  using assms by auto

theorem square-test-correct: square-test n ⟷ is-square n
proof cases
  assume is-square n
  hence rhs: ∃ q. q2 = n by (auto elim: is-nth-powerE)
  note sq-mod = square-mod-existence[OF this]
  have q64-member: n mod 64 ∈ q64 using sq-mod[of 64]
    unfolding q64-infinite-def image-def by simp
  let ?r = n mod 45045
  have 11 dvd (45045::nat) 63 dvd (45045::nat) 65 dvd (45045::nat) by force+
  then have mod-45045: ?r mod 11 = n mod 11 ?r mod 63 = n mod 63 ?r mod
65 = n mod 65
    using mod-mod-cancel[of - 45045 n] by presburger+
  then have ?r mod 11 ∈ q11 ?r mod 63 ∈ q63 ?r mod 65 ∈ q65
    using sq-mod[of 11] sq-mod[of 63] sq-mod[of 65]
    unfolding q11-infinite-def q63-infinite-def q65-infinite-def image-def mod-45045
    by fast+
  then show ?thesis unfolding square-test-def Let-def using q64-member rhs by
auto
next
  assume not-rhs: ¬is-square n
  hence ∄ q. q2 = n by auto
  then have (Discrete.sqrt n)2 ≠ n by simp
  then show ?thesis unfolding square-test-def by (auto simp: is-nth-power-def)
qed

```


definition *get-nat-sqrt* :: *nat* \Rightarrow *nat option*
where *get-nat-sqrt* *n* = (*if is-square n then Some (Discrete.sqrt n) else None*)

lemma *get-nat-sqrt-code* [*code*]:
get-nat-sqrt *n* =
 (*if IArray.sub q64-array (n mod 64) \wedge (let r = n mod 45045 in*
 IArray.sub q63-array (r mod 63) \wedge
 IArray.sub q65-array (r mod 65) \wedge
 IArray.sub q11-array (r mod 11)) then
 (*let x = Discrete.sqrt n in if x² = n then Some x else None*) *else None*)
unfolding *get-nat-sqrt-def square-test-correct* [*symmetric*] *square-test-def*
using *in-q11-code* [*symmetric*] *in-q63-code* [*symmetric*]
 in-q64-code [*symmetric*] *in-q65-code* [*symmetric*]
by (*auto split: if-splits simp: Let-def*)

end

2 Pell's equation

theory *Pell*

imports

Complex-Main

HOL-Computational-Algebra.Computational-Algebra

begin

Pell's equation has the general form $x^2 = 1 + Dy^2$ where $D \in \mathbb{N}$ is a parameter and x, y are \mathbb{Z} -valued variables. As we will see, that case where D is a perfect square is trivial and therefore uninteresting; we will therefore assume that D is not a perfect square for the most part.

Furthermore, it is obvious that the solutions to Pell's equation are symmetric around the origin in the sense that (x, y) is a solution iff $(\pm x, \pm y)$ is a solution. We will therefore mostly look at solutions (x, y) where both x and y are non-negative, since the remaining solutions are a trivial consequence of these.

Information on the material treated in this formalisation can be found in many textbooks and lecture notes, e. g. [3, 1].

2.1 Preliminary facts

lemma *gcd-int-nonpos-iff* [*simp*]: $\text{gcd } x (y :: \text{int}) \leq 0 \iff x = 0 \wedge y = 0$

proof

assume $\text{gcd } x y \leq 0$

with *gcd-ge-0-int*[*of x y*] **have** $\text{gcd } x y = 0$ **by** *linarith*

thus $x = 0 \wedge y = 0$ **by** *auto*

qed *auto*

lemma *minus-in-Ints-iff* [*simp*]:

$-x \in \mathbb{Z} \longleftrightarrow x \in \mathbb{Z}$
using *Ints-minus[of x] Ints-minus[of -x]* **by** *auto*

A (positive) square root of a natural number is either a natural number or irrational.

lemma *nonneg-sqrt-nat-or-irrat*:

assumes $x^2 = \text{real } a$ **and** $x \geq 0$

shows $x \in \mathbb{N} \vee x \notin \mathbb{Q}$

proof *safe*

assume $x \notin \mathbb{N}$ **and** $x \in \mathbb{Q}$

from *Rats-abs-nat-div-natE[OF this(2)]*

obtain $p \ q :: \text{nat}$ **where** $q \neq 0$ **and** $\text{abs } x = p / q$ **and** *coprime*:
coprime p q .

with $\langle x \geq 0 \rangle$ **have** $x = p / q$

by *simp*

with *assms* **have** $\text{real } (q^2) * \text{real } a = \text{real } (p^2)$

by (*simp add: field-simps*)

also **have** $\text{real } (q^2) * \text{real } a = \text{real } (q^2 * a)$

by *simp*

finally **have** $p^2 = q^2 * a$

by (*subst (asm) of-nat-eq-iff*) *auto*

hence $q^2 \text{ dvd } p^2$

by *simp*

hence $q \text{ dvd } p$

by *simp*

with *coprime* **have** $q = 1$

by *auto*

with x **and** $\langle x \notin \mathbb{N} \rangle$ **show** *False*

by *simp*

qed

A square root of a natural number is either an integer or irrational.

corollary *sqrt-nat-or-irrat*:

assumes $x^2 = \text{real } a$

shows $x \in \mathbb{Z} \vee x \notin \mathbb{Q}$

proof (*cases x ≥ 0*)

case *True*

with *nonneg-sqrt-nat-or-irrat[OF assms this]*

show *?thesis* **by** (*auto simp: Nats-altdef2*)

next

case *False*

from *assms* **have** $(-x)^2 = \text{real } a$

by *simp*

moreover **from** *False* **have** $-x \geq 0$

by *simp*

ultimately **have** $-x \in \mathbb{N} \vee -x \notin \mathbb{Q}$

by (*rule nonneg-sqrt-nat-or-irrat*)

thus *?thesis*

by (*auto simp: Nats-altdef2*)

qed

corollary *sqrt-nat-or-irrat'*:

sqrt (real a) ∈ ℕ ∨ sqrt (real a) ∉ ℚ
using *nonneg-sqrt-nat-or-irrat[of sqrt a a]* **by** *auto*

The square root of a natural number n is again a natural number iff n is a perfect square.

corollary *sqrt-nat-iff-is-square*:

sqrt (real n) ∈ ℕ ↔ is-square n

proof

assume *sqrt (real n) ∈ ℕ*

then obtain k **where** *sqrt (real n) = real k* **by** (*auto elim!: Nats-cases*)

hence *sqrt (real n) ^ 2 = real (k ^ 2)* **by** (*simp only: of-nat-power*)

also have *sqrt (real n) ^ 2 = real n* **by** *simp*

finally have $n = k ^ 2$ **by** (*simp only: of-nat-eq-iff*)

thus *is-square n* **by** *blast*

qed (*auto elim!: is-nth-powerE*)

corollary *irrat-sqrt-nonsquare*: $\neg is-square n \implies sqrt (real n) \notin \mathbb{Q}$

using *sqrt-nat-or-irrat'[of n]* **by** (*auto simp: sqrt-nat-iff-is-square*)

2.2 The case of a perfect square

As we have noted, the case where D is a perfect square is trivial: In fact, we will show that the only solutions in this case are the trivial solutions $(x, y) = (\pm 1, 0)$ if D is a non-zero perfect square, or $(\pm 1, y)$ for arbitrary $y \in \mathbb{Z}$ if $D = 0$.

context

fixes $D :: nat$

assumes *square-D: is-square D*

begin

lemma *pell-square-solution-nat-aux*:

fixes $x y :: nat$

assumes $D > 0$ **and** $x ^ 2 = 1 + D * y ^ 2$

shows $(x, y) = (1, 0)$

proof –

from *assms* **have** $x > 0$ **by** (*auto intro!: Nat.gr0I*)

from *square-D* **obtain** d **where** [*simp*]: $D = d^2$

by (*auto elim: is-nth-powerE*)

have $int x ^ 2 = int (x ^ 2)$ **by** *simp*

also note *assms(2)*

also have $int (1 + D * y ^ 2) = 1 + int D * int y ^ 2$ **by** *simp*

finally have $(int x + int d * int y) * (int x - int d * int y) = 1$

by (*simp add: algebra-simps power2-eq-square*)

hence $*$: $int x + int d * int y = 1 \wedge int x - int d * int y = 1$

using $x > 0$ **by** (*subst (asm) pos-zmult-eq-1-iff*) (*auto intro: add-pos-nonneg*)

from * **have** [simp]: $x = 1$ **by** simp
moreover from * **and** *assms(1)* **have** $y = 0$ **by** auto
ultimately show ?thesis **by** simp
qed

lemma *pell-square-solution-int-aux*:

fixes $x\ y :: \text{int}$
assumes $D > 0$ **and** $x^2 = 1 + D * y^2$
shows $x \in \{-1, 1\} \wedge y = 0$
proof –
define $x'\ y'$ **where** $x' = \text{nat } |x|$ **and** $y' = \text{nat } |y|$
have $x = \text{sgn } x * x'$ **and** $y = \text{sgn } y * y'$
by (auto simp: sgn-if x'-def y'-def)
have zero-iff: $x = 0 \iff x' = 0$ $y = 0 \iff y' = 0$
by (auto simp: x'-def y'-def)
note *assms(2)*
also have $x^2 = \text{int } (x'^2)$
by (subst x) (auto simp: sgn-if zero-iff)
also have $1 + D * y^2 = \text{int } (1 + D * y'^2)$
by (subst y) (auto simp: sgn-if zero-iff)
also note of-nat-eq-iff
finally have $x'^2 = 1 + D * y'^2$.
from $\langle D > 0 \rangle$ **and this have** $(x', y') = (1, 0)$
by (rule pell-square-solution-nat-aux)
thus ?thesis **by** (auto simp: x'-def y'-def)
qed

lemma *pell-square-solution-nat-iff*:

fixes $x\ y :: \text{nat}$
shows $x^2 = 1 + D * y^2 \iff x = 1 \wedge (D = 0 \vee y = 0)$
using *pell-square-solution-nat-aux*[of $x\ y$] **by** (cases $D = 0$) auto

lemma *pell-square-solution-int-iff*:

fixes $x\ y :: \text{int}$
shows $x^2 = 1 + D * y^2 \iff x \in \{-1, 1\} \wedge (D = 0 \vee y = 0)$
using *pell-square-solution-int-aux*[of $x\ y$] **by** (cases $D = 0$) (auto simp: power2-eq-1-iff)

end

2.3 Existence of a non-trivial solution

Let us now turn to the case where D is not a perfect square.

We first show that Pell's equation always has at least one non-trivial solution (apart from the trivial solution $(1, 0)$). For this, we first need a lemma about the existence of rational approximations of real numbers.

The following lemma states that for any positive integer s and real number x , we can find a rational approximation t / u to x with an error of most $1 / (u * s)$ where the denominator u is at most s .

lemma *pell-approximation-lemma*:
fixes $s :: \text{nat}$ **and** $x :: \text{real}$
assumes $s : s > 0$
shows $\exists u :: \text{nat}. \exists t :: \text{int}. u > 0 \wedge \text{coprime } u \ t \wedge 1 / s \in \{|t - u * x| < .1 / u\}$
proof –
define f **where** $f = (\lambda u. \lceil u * x \rceil)$
define $g :: \text{nat} \Rightarrow \text{int}$ **where** $g = (\lambda u. \lfloor \text{frac } (u * x) * s \rfloor)$
{
 fix $u :: \text{nat}$ **assume** $u : u \leq s$
 hence $\text{frac } (u * x) * \text{real } s < 1 * \text{real } s$
 using s **by** (*intro mult-strict-right-mono*) (*auto simp: frac-lt-1*)
 hence $g \ u < \text{int } s$ **by** (*auto simp: floor-less-iff g-def*)
}
hence $g \ \{..s\} \subseteq \{0..<s\}$
 by (*auto simp: g-def floor-less-iff*)
hence $\text{card } (g \ \{..s\}) \leq \text{card } \{0..<\text{int } s\}$
 by (*intro card-mono*) *auto*
also have $\dots < \text{card } \{..s\}$ **by** *simp*
finally have $\neg \text{inj-on } g \ \{..s\}$ **by** (*rule pigeonhole*)
then obtain $a \ b$ **where** $ab : a \leq s \ b \leq s \ a \neq b \ g \ a = g \ b$
 by (*auto simp: inj-on-def*)
define $u1$ **and** $u2$ **where** $u1 = \max \ a \ b$ **and** $u2 = \min \ a \ b$
have $u12 : u1 \leq s \ u2 \leq s \ u2 < u1 \ g \ u1 = g \ u2$
 using ab **by** (*auto simp: u1-def u2-def*)

define $u \ t$ **where** $u = u1 - u2$ **and** $t = \lfloor u1 * x \rfloor - \lfloor u2 * x \rfloor$
have $u : u > 0 \ |u| \leq s$
 using $u12$ **by** (*simp-all add: u-def*)

from $\langle g \ u1 = g \ u2 \rangle$ **have** $|\text{frac } (u2 * x) * s - \text{frac } (u1 * x) * s| < 1$
 unfolding $g\text{-def}$ **by** *linarith*
also have $|\text{frac } (u2 * x) * s - \text{frac } (u1 * x) * s| =$
 $|\text{real } s| * |\text{frac } (u2 * x) - \text{frac } (u1 * x)|$
 by (*subst abs-mult [symmetric]*) (*simp add: algebra-simps*)
finally have $|t - u * x| * s < 1$ **using** $\langle u1 > u2 \rangle$
 by (*simp add: g-def u-def t-def frac-def algebra-simps of-nat-diff*)
with $\langle s > 0 \rangle$ **have** *less*: $|t - u * x| < 1 / s$ **by** (*simp add: divide-simps*)

define d **where** $d = \text{gcd } (\text{nat } |t|) \ u$
define $t' :: \text{int}$ **and** $u' :: \text{nat}$ **where** $t' = t \ \text{div } d$ **and** $u' = u \ \text{div } d$
from u **have** $d \neq 0$
 by (*intro notI*) (*auto simp: d-def*)
have $\text{int } (\text{gcd } (\text{nat } |t|) \ u) = \text{gcd } |t| \ (\text{int } u)$
 by *simp*
hence $t' \cdot u' : t = t' * d \ u = u' * d$
 by (*auto simp: t'-def u'-def d-def nat-dvd-iff*)

from $\langle d \neq 0 \rangle$ **have** $|t' - u' * x| * 1 \leq |t' - u' * x| * |\text{real } d|$
 by (*intro mult-left-mono*) *auto*

also have $\dots = |t - u * x|$ **by** (*subst abs-mult [symmetric]*) (*simp add: algebra-simps t'-u'*)
also note *less*
finally have $|t' - u' * x| < 1 / s$ **by** *simp*
moreover {
 from $\langle s > 0 \rangle$ **and** u **have** $1 / s \leq 1 / u$
 by (*simp add: divide-simps u-def*)
 also have $\dots = 1 / u' / d$ **by** (*simp add: t'-u' divide-simps*)
 also have $\dots \leq 1 / u' / 1$ **using** $\langle d \neq 0 \rangle$ **by** (*intro divide-left-mono*) *auto*
 finally have $1 / s \leq 1 / u'$ **by** *simp*
}
ultimately have $1 / s \in \{|t' - u' * x| < ..1 / u'\}$ **by** *auto*
moreover from $\langle u > 0 \rangle$ **have** $u' > 0$ **by** (*auto simp: t'-u'*)
moreover {
 have $\text{gcd } u \ t = \text{gcd } t' \ u' * \text{int } d$
 by (*simp add: t'-u' gcd-mult-right gcd.commute*)
 also have $\text{int } d = \text{gcd } u \ t$
 by (*simp add: d-def gcd.commute*)
 finally have $\text{gcd } u' \ t' = 1$ **using** u **by** (*simp add: gcd.commute*)
}
ultimately show *?thesis* **by** *blast*
qed

As a simple corollary of this, we can show that for irrational x , there is an infinite number of rational approximations t / u to x whose error is less than $1 / u^2$.

corollary *pell-approximation-corollary*:

fixes $x :: \text{real}$
assumes $x \notin \mathbb{Q}$
shows *infinite* $\{(t :: \text{int}, u :: \text{nat}). u > 0 \wedge \text{coprime } u \ t \wedge |t - u * x| < 1 / u\}$
 (*is infinite ?A*)
proof
 assume *fin: finite ?A*
 let $?f = \lambda(t :: \text{int}, u :: \text{nat}). |t - u * x|$
 from *fin* **have** *fin'*: *finite* (*insert 1 (?f ' ?A)*) **by** *blast*
 have $\text{Min } (\text{insert } 1 \ (?f \ ' \ ?A)) > 0$
 proof (*subst Min-gr-iff*)
 have $a \neq b * x$ **if** $b > 0$ **for** $a :: \text{int}$ **and** $b :: \text{nat}$
 proof
 assume $a = b * x$
 with $\langle b > 0 \rangle$ **have** $x = a / b$ **by** (*simp add: field-simps*)
 with $\langle x \notin \mathbb{Q} \rangle$ **and** $\langle b > 0 \rangle$ **show** *False* **by** (*auto simp: Rats-eq-int-div-nat*)
 qed
 thus $\forall x \in \text{insert } 1 \ (?f \ ' \ ?A). x > 0$ **by** *auto*
 qed (*insert fin', simp-all*)
 also note *real-arch-inverse*
 finally obtain $M :: \text{nat}$ **where** $M: M \neq 0$ *inverse* $M < \text{Min } (\text{insert } 1 \ (?f \ ' \ ?A))$
 by *blast*
 hence $M > 0$ **by** *simp*

```

from pell-approximation-lemma[OF this, of x] obtain  $u :: \text{nat}$  and  $t :: \text{int}$ 
  where  $ut: u > 0$  coprime u t  $1 / \text{real } M \in \{?f (t, u) <..1 / u\}$  by auto
from  $ut$  have  $?f (t, u) < 1 / \text{real } M$  by simp
also from  $M$  have  $\dots < \text{Min} (\text{insert } 1 (?f ' ?A))$ 
  by (simp add: divide-simps)
also from  $ut$  have  $\text{Min} (\text{insert } 1 (?f ' ?A)) \leq ?f (t, u)$ 
  by (intro Min.coboundedI fin') auto
finally show False by simp
qed

```

```

locale pell =
  fixes  $D :: \text{nat}$ 
  assumes nonsquare-D:  $\neg$ is-square D
begin

```

```

lemma D-gt-1:  $D > 1$ 
proof -
  from nonsquare-D have  $D \neq 0$   $D \neq 1$  by (auto intro!: Nat.gr0I)
  thus ?thesis by simp
qed

```

```

lemma D-pos:  $D > 0$ 
  using nonsquare-D by (intro Nat.gr0I) auto

```

With the above corollary, we can show the existence of a non-trivial solution. We restrict our attention to solutions (x, y) where both x and y are non-negative.

```

theorem pell-solution-exists:  $\exists (x::\text{nat}) (y::\text{nat}). y \neq 0 \wedge x^2 = 1 + D * y^2$ 
proof -
  define  $S$  where  $S = \{(t :: \text{int}, u :: \text{nat}). u > 0 \wedge \text{coprime } u \ t \wedge |t - u * \text{sqrt } D| < 1 / u\}$ 
  let  $?f = \lambda(t :: \text{int}, u :: \text{nat}). t^2 - u^2 * D$ 
  define  $M$  where  $M = \lfloor 1 + 2 * \text{sqrt } D \rfloor$ 
  have infinite:  $\neg$ finite S unfolding S-def
    by (intro pell-approximation-corollary irrat-sqrt-nonsquare nonsquare-D)

  have subset:  $?f ' S \subseteq \{-M..M\}$ 
proof safe
  fix  $u :: \text{nat}$  and  $t :: \text{int}$ 
  assume  $tu: (t, u) \in S$ 
  from  $tu$  have [simp]:  $u > 0$  by (auto simp: S-def)
  have  $|t + u * \text{sqrt } D| = |t - u * \text{sqrt } D + 2 * u * \text{sqrt } D|$  by simp
  also have  $\dots \leq |t - u * \text{sqrt } D| + |2 * u * \text{sqrt } D|$ 
    by (rule abs-triangle-ineq)
  also have  $|2 * u * \text{sqrt } D| = 2 * u * \text{sqrt } D$  by simp
  also have  $|t - u * \text{sqrt } D| \leq 1 / u$ 
    using  $tu$  by (simp add: S-def)

```

finally have $le: |t + u * \text{sqrt } D| \leq 1 / u + 2 * u * \text{sqrt } D$ **by** *simp*

have $|t^2 - u^2 * D| = |t - u * \text{sqrt } D| * |t + u * \text{sqrt } D|$
by (*subst abs-mult [symmetric]*) (*simp add: algebra-simps power2-eq-square*)

also have $\dots \leq 1 / u * (1 / u + 2 * u * \text{sqrt } D)$
using *tu* **by** (*intro mult-mono le*) (*auto simp: S-def*)

also have $\dots = 1 / \text{real } u ^ 2 + 2 * \text{sqrt } D$
by (*simp add: algebra-simps power2-eq-square*)

also from $\langle u > 0 \rangle$ **have** $\text{real } u \geq 1$ **by** *linarith*

hence $1 / \text{real } u ^ 2 \leq 1 / 1 ^ 2$
by (*intro divide-left-mono power-mono*) *auto*

finally have $|t^2 - u^2 * D| \leq 1 + 2 * \text{sqrt } D$ **by** *simp*

hence $t^2 - u^2 * D \geq -M$ $t^2 - u^2 * D \leq M$ **unfolding** *M-def* **by** *linarith+*

thus $t^2 - u^2 * D \in \{-M..M\}$ **by** *simp*

qed

hence *fin*: *finite* (*?f ' S*) **by** (*rule finite-subset*) *auto*

from *pigeonhole-infinite*[*OF infinite fin*]

obtain *z* **where** $z: z \in S$ *infinite* $\{z' \in S. ?f z' = ?f z\}$ **by** *blast*

define *k* **where** $k = ?f z$

with *subset* **and** *z* **have** $k: k \in \{-M..M\}$ *infinite* $\{z \in S. ?f z = k\}$
by (*auto simp: k-def*)

have *k-nz*: $k \neq 0$

proof

assume [*simp*]: $k = 0$

note *k(2)*

also have $?f z \neq 0$ **if** $z \in S$ **for** *z*

proof

assume *: $?f z = 0$

obtain *t u* **where** [*simp*]: $z = (t, u)$ **by** (*cases z*)

from * **have** $t ^ 2 = \text{int } u ^ 2 * \text{int } D$ **by** *simp*

hence $\text{int } u ^ 2 \text{ dvd } t ^ 2$ **by** *simp*

hence $\text{int } u \text{ dvd } t$ **by** *simp*

then obtain *k* **where** [*simp*]: $t = \text{int } u * k$ **by** (*auto elim!: dvdE*)

from * **and** $\langle z \in S \rangle$ **have** $k ^ 2 = \text{int } D$
by (*auto simp: power-mult-distrib S-def*)

also have $k ^ 2 = \text{int } (\text{nat } |k| ^ 2)$ **by** *simp*

finally have $D = \text{nat } |k| ^ 2$ **by** (*simp only: of-nat-eq-iff*)

hence *is-square* *D* **by** *auto*

with *nonsquare-D* **show** *False* **by** *contradiction*

qed

hence $\{z \in S. ?f z = k\} = \{\}$ **by** *auto*

finally show *False* **by** *simp*

qed

let *?h* = $\lambda(t :: \text{int}, u :: \text{nat}). (t \bmod (\text{abs } k), u \bmod (\text{abs } k))$

have *?h* ' $\{z \in S. ?f z = k\} \subseteq \{0..<\text{abs } k\} \times \{0..<\text{abs } k\}$
using *k-nz* **by** (*auto simp: case-prod-unfold*)

hence *finite* ($?h \{z \in S. ?f z = k\}$) **by** (*rule finite-subset*) *auto*
from *pigeonhole-infinite*[*OF k(2) this*] **obtain** z'
where $z': z' \in S ?f z' = k$ *infinite* $\{z'' \in \{z \in S. ?f z = k\}. ?h z'' = ?h z'\}$
by *blast*
define $l1$ **and** $l2$ **where** $l1 = fst (?h z')$ **and** $l2 = snd (?h z')$
define S' **where** $S' = \{(t,u) \in S. ?f (t,u) = k \wedge t \bmod abs k = l1 \wedge u \bmod abs k = l2\}$
note $z'(3)$
also have $\{z'' \in \{z \in S. ?f z = k\}. ?h z'' = ?h z'\} = S'$
by (*auto simp: l1-def l2-def case-prod-unfold S'-def*)
finally have *infinite: infinite S' .*

from $z'(1)$ **and** $k-nz$ **have** $l12: l1 \in \{0..<abs k\} l2 \in \{0..<abs k\}$
by (*auto simp: l1-def l2-def case-prod-unfold*)

from *infinite-arbitrarily-large*[*OF infinite*]
obtain X **where** $X: finite X \ card X = 2 \ X \subseteq S'$ **by** *blast*
from *finite-distinct-list*[*OF this(1)*] **obtain** xs **where** $xs: set xs = X$ *distinct xs*
by *blast*
with X **have** $length xs = 2$ **using** *distinct-card*[*of xs*] **by** *simp*
then obtain $z1 z2$ **where** [*simp*]: $xs = [z1, z2]$
by (*auto simp: length-Suc-conv eval-nat-numeral*)
from $X xs$ **have** $S': z1 \in S' z2 \in S'$ **and** *neg: z1 ≠ z2* **by** *auto*
define $t1 u1 t2 u2$ **where** $t1 = fst z1$ **and** $u1 = snd z1$ **and** $t2 = fst z2$ **and**
 $u2 = snd z2$
have [*simp*]: $z1 = (t1, u1) z2 = (t2, u2)$
by (*simp-all add: t1-def u1-def t2-def u2-def*)

from S' **have** * [*simp*]: $t1 \bmod abs k = l1 \ t2 \bmod abs k = l1 \ u1 \bmod abs k = l2$
 $u2 \bmod abs k = l2$
by (*simp-all add: S'-def*)
define x **where** $x = (t1 * t2 - D * u1 * u2) \ div k$
define y **where** $y = (t1 * u2 - t2 * u1) \ div k$

from S' **have** $(t1^2 - u1^2 * D) = k \ (t2^2 - u2^2 * D) = k$
by (*auto simp: S'-def*)
hence $(t1^2 - u1^2 * D) * (t2^2 - u2^2 * D) = k ^ 2$
unfolding *power2-eq-square* **by** *simp*
also have $(t1^2 - u1^2 * D) * (t2^2 - u2^2 * D) =$
 $(t1 * t2 - D * u1 * u2) ^ 2 - D * (t1 * u2 - t2 * u1) ^ 2$
by (*simp add: power2-eq-square algebra-simps*)
finally have *eq: (t1 * t2 - D * u1 * u2)^2 - D * (t1 * u2 - t2 * u1)^2 = k^2 .*

have $(t1 * u2 - t2 * u1) \bmod abs k = (l1 * l2 - l1 * l2) \bmod abs k$
using $l12$ **by** (*intro mod-diff-cong mod-mult-cong*) (*auto simp: mod-pos-pos-trivial*)
hence $dvd1: k \ dvd \ t1 * u2 - t2 * u1$ **by** (*simp add: mod-eq-0-iff-dvd*)

have $k^2 \ dvd \ k^2 + D * (t1 * u2 - t2 * u1)^2$
using $dvd1$ **by** (*intro dvd-add*) *auto*

also from *eq* **have** $\dots = (t1 * t2 - D * u1 * u2)^2$
by (*simp add: algebra-simps*)
finally have *dvd2*: $k \text{ dvd } t1 * t2 - D * u1 * u2$
by *simp*

note *eq*
also from *dvd2* **have** $t1 * t2 - D * u1 * u2 = k * x$
by (*simp add: x-def*)
also from *dvd1* **have** $t1 * u2 - t2 * u1 = k * y$
by (*simp add: y-def*)
also have $(k * x)^2 - D * (k * y)^2 = k^2 * (x^2 - D * y^2)$
by (*simp add: power-mult-distrib algebra-simps*)
finally have *eq'*: $x^2 - D * y^2 = 1$
using *k-nz* **by** *simp*
hence $x^2 = 1 + D * y^2$ **by** *simp*
also have $x^2 = \text{int } (\text{nat } |x| \wedge 2)$ **by** *simp*
also have $1 + D * y^2 = \text{int } (1 + D * \text{nat } |y| \wedge 2)$ **by** *simp*
also note *of-nat-eq-iff*
finally have *eq''*: $(\text{nat } |x|)^2 = 1 + D * (\text{nat } |y|)^2$.

have $t1 * u2 \neq t2 * u1$
proof
assume *: $t1 * u2 = t2 * u1$
hence $|t1| * |u2| = |t2| * |u1|$ **by** (*simp only: abs-mult [symmetric]*)
moreover from *S'* **have** *coprime u1 t1 coprime u2 t2*
by (*auto simp: S'-def S-def*)
ultimately have *eq*: $|t1| = |t2| \wedge u1 = u2$
by (*subst (asm) coprime-crossproduct-int*) (*auto simp: S'-def S-def gcd.commute coprime-commute*)
moreover from *S'* **have** $u1 \neq 0 \ u2 \neq 0$ **by** (*auto simp: S'-def S-def*)
ultimately have $t1 = t2$ **using** * **by** *auto*
with *eq* **and** *neq* **show** *False* **by** *auto*
qed
with *dvd1* **have** $y \neq 0$
by (*auto simp add: y-def dvd-div-eq-0-iff*)
hence $\text{nat } |y| \neq 0$ **by** *auto*
with *eq''* **show** $\exists x \ y. y \neq 0 \wedge x^2 = 1 + D * y^2$ **by** *blast*
qed

2.4 Definition of solutions

We define some abbreviations for the concepts of a solution and a non-trivial solution.

definition *solution* :: $('a \times 'a :: \text{comm-semiring-1}) \Rightarrow \text{bool}$ **where**
solution = $(\lambda(a, b). a^2 = 1 + \text{of-nat } D * b^2)$

definition *nontriv-solution* :: $('a \times 'a :: \text{comm-semiring-1}) \Rightarrow \text{bool}$ **where**
nontriv-solution = $(\lambda(a, b). (a, b) \neq (1, 0) \wedge a^2 = 1 + \text{of-nat } D * b^2)$

lemma *nontriv-solution-altdef*: $\text{nontriv-solution } z \longleftrightarrow \text{solution } z \wedge z \neq (1, 0)$
by (*auto simp: solution-def nontriv-solution-def*)

lemma *solution-trivial-nat* [*simp, intro*]: $\text{solution } (\text{Suc } 0, 0)$
by (*simp add: solution-def*)

lemma *solution-trivial* [*simp, intro*]: $\text{solution } (1, 0)$
by (*simp add: solution-def*)

lemma *solution-uminus-left* [*simp*]: $\text{solution } (-x, y :: 'a :: \text{comm-ring-1}) \longleftrightarrow \text{solution } (x, y)$
by (*simp add: solution-def*)

lemma *solution-uminus-right* [*simp*]: $\text{solution } (x, -y :: 'a :: \text{comm-ring-1}) \longleftrightarrow \text{solution } (x, y)$
by (*simp add: solution-def*)

lemma *solution-0-snd-nat-iff* [*simp*]: $\text{solution } (a :: \text{nat}, 0) \longleftrightarrow a = 1$
by (*auto simp: solution-def*)

lemma *solution-0-snd-iff* [*simp*]: $\text{solution } (a :: 'a :: \text{idom}, 0) \longleftrightarrow a \in \{1, -1\}$
by (*auto simp: solution-def power2-eq-1-iff*)

lemma *no-solution-0-fst-nat* [*simp*]: $\neg \text{solution } (0, b :: \text{nat})$
by (*auto simp: solution-def*)

lemma *no-solution-0-fst-int* [*simp*]: $\neg \text{solution } (0, b :: \text{int})$
proof –
have $1 + \text{int } D * b^2 > 0$ **by** (*intro add-pos-nonneg auto*)
thus *?thesis* **by** (*auto simp add: solution-def*)
qed

lemma *solution-of-nat-of-nat* [*simp*]:
 $\text{solution } (\text{of-nat } a, \text{of-nat } b :: 'a :: \{\text{comm-ring-1}, \text{ring-char-0}\}) \longleftrightarrow \text{solution } (a, b)$
by (*simp only: solution-def prod.case of-nat-power [symmetric]*
of-nat-1 [symmetric, where ?'a = 'a] of-nat-add [symmetric]
of-nat-mult [symmetric] of-nat-eq-iff of-nat-id)

lemma *solution-of-nat-of-nat'* [*simp*]:
 $\text{solution } (\text{case } z \text{ of } (a, b) \Rightarrow (\text{of-nat } a, \text{of-nat } b :: 'a :: \{\text{comm-ring-1}, \text{ring-char-0}\}))$
 \longleftrightarrow
 $\text{solution } z$
by (*auto simp: case-prod-unfold*)

lemma *solution-nat-abs-nat-abs* [*simp*]:
 $\text{solution } (\text{nat } |x|, \text{nat } |y|) \longleftrightarrow \text{solution } (x, y)$
proof –
define x' **and** y' **where** $x' = \text{nat } |x|$ **and** $y' = \text{nat } |y|$

have $x: x = \text{sgn } x * x'$ **and** $y: y = \text{sgn } y * y'$
by (*auto simp: x'-def y'-def sgn-if*)
have [*simp*]: $x = 0 \longleftrightarrow x' = 0$ $y = 0 \longleftrightarrow y' = 0$
by (*auto simp: x'-def y'-def*)
show *solution* (x', y') \longleftrightarrow *solution* (x, y)
by (*subst x, subst y*) (*auto simp: sgn-if*)
qed

lemma *nontriv-solution-of-nat-of-nat* [*simp*]:
nontriv-solution (*of-nat a, of-nat b* :: ' a ' :: {*comm-ring-1, ring-char-0*}) \longleftrightarrow
nontriv-solution (a, b)
by (*auto simp: nontriv-solution-altdef*)

lemma *nontriv-solution-of-nat-of-nat'* [*simp*]:
nontriv-solution (*case z of* $(a, b) \Rightarrow$ (*of-nat a, of-nat b* :: ' a ' :: {*comm-ring-1, ring-char-0*})) \longleftrightarrow
nontriv-solution z
by (*auto simp: case-prod-unfold*)

lemma *nontriv-solution-imp-solution* [*dest*]: *nontriv-solution* $z \implies$ *solution* z
by (*auto simp: nontriv-solution-altdef*)

2.5 The Pell valuation function

Solutions (x, y) have an interesting correspondence to the ring $\mathbb{Z}[\sqrt{D}]$ via the map $(x, y) \mapsto x + y\sqrt{D}$. We call this map the *Pell valuation function*. It is obvious that this map is injective, since \sqrt{D} is irrational.

definition *pell-valuation* :: *int* \times *int* \Rightarrow *real* **where**
pell-valuation = $(\lambda(a, b). a + b * \text{sqrt } D)$

lemma *pell-valuation-nonneg* [*simp*]: *fst z* $\geq 0 \implies$ *snd z* $\geq 0 \implies$ *pell-valuation z* ≥ 0
by (*auto simp: pell-valuation-def case-prod-unfold*)

lemma *pell-valuation-uminus-uminus* [*simp*]: *pell-valuation* $(-x, -y) = -$ *pell-valuation* (x, y)
by (*simp add: pell-valuation-def*)

lemma *pell-valuation-eq-iff* [*simp*]:
pell-valuation z1 = *pell-valuation z2* \longleftrightarrow $z1 = z2$

proof

assume *: *pell-valuation z1* = *pell-valuation z2*

obtain a b **where** [*simp*]: $z1 = (a, b)$ **by** (*cases z1*)

obtain u v **where** [*simp*]: $z2 = (u, v)$ **by** (*cases z2*)

have $b = v$

proof (*rule ccontr*)

assume $b \neq v$

with * **have** $\text{sqrt } D = (u - a) / (b - v)$

by (*simp add: field-simps pell-valuation-def*)

also have $\dots \in \mathbb{Q}$ by *auto*
 finally show *False* using *irrat-sqrt-nonsquare nonsquare-D* by *blast*
 qed
 moreover from *this* and $*$ have $a = u$
 by (*simp add: pell-valuation-def*)
 ultimately show $z1 = z2$ by *simp*
 qed *auto*

2.6 Linear ordering of solutions

Next, we show that solutions are linearly ordered w. r. t. the pointwise order on products. This means that for two different solutions (a, b) and (x, y) , we always either have $a < x$ and $b < y$ or $a > x$ and $b > y$.

lemma *solutions-linorder*:

fixes $a b x y :: \text{nat}$
 assumes *solution* (a, b) *solution* (x, y)
 shows $a \leq x \wedge b \leq y \vee a \geq x \wedge b \geq y$
proof –
 have $b \leq y$ if $a \leq x$ *solution* (a, b) *solution* (x, y) for $a b x y :: \text{nat}$
proof –
 from *that* have $a^2 \leq x^2$ by (*intro power-mono*) *auto*
 with *that* and *D-gt-1* have $b^2 \leq y^2$
 by (*simp add: solution-def*)
 thus $b \leq y$
 by (*simp add: power2-nat-le-eq-le*)
 qed
 from *this*[of $a x b y$] and *this*[of $x a y b$] and *assms* show *?thesis*
 by (*cases a ≤ x*) *auto*
 qed

lemma *solutions-linorder-strict*:

fixes $a b x y :: \text{nat}$
 assumes *solution* (a, b) *solution* (x, y)
 shows $(a, b) = (x, y) \vee a < x \wedge b < y \vee a > x \wedge b > y$
proof –
 have $b = y$ if $a = x$
 using *that* *assms* and *D-gt-1* by (*simp add: solution-def*)
 moreover have $a = x$ if $b = y$
proof –
 from *that* and *assms* have $a^2 = \text{Suc } (D * y^2)$
 by (*simp add: solution-def*)
 also from *that* and *assms* have $\dots = x^2$
 by (*simp add: solution-def*)
 finally show $a = x$ by *simp*
 qed
 ultimately have [*simp*]: $a = x \longleftrightarrow b = y$..
 show *?thesis* using *solutions-linorder*[*OF assms*]
 by (*cases a x* rule: *linorder-cases*; *cases b y* rule: *linorder-cases*) *simp-all*
 qed

```

lemma solutions-le-iff-pell-valuation-le:
  fixes a b x y :: nat
  assumes solution (a, b) solution (x, y)
  shows a ≤ x ∧ b ≤ y ⟷ pell-valuation (a, b) ≤ pell-valuation (x, y)
proof
  assume a ≤ x ∧ b ≤ y
  thus pell-valuation (a, b) ≤ pell-valuation (x, y)
    unfolding pell-valuation-def prod.case using D-gt-1
    by (intro add-mono mult-right-mono) auto
next
  assume *: pell-valuation (a, b) ≤ pell-valuation (x, y)
  from assms have a ≤ x ∧ b ≤ y ∨ x ≤ a ∧ y ≤ b
    by (rule solutions-linorder)
  thus a ≤ x ∧ b ≤ y
proof
  assume x ≤ a ∧ y ≤ b
  hence pell-valuation (a, b) ≥ pell-valuation (x, y)
    unfolding pell-valuation-def prod.case using D-gt-1
    by (intro add-mono mult-right-mono) auto
  with * have pell-valuation (a, b) = pell-valuation (x, y) by linarith
  hence (a, b) = (x, y) by simp
  thus a ≤ x ∧ b ≤ y by simp
qed auto
qed

lemma solutions-less-iff-pell-valuation-less:
  fixes a b x y :: nat
  assumes solution (a, b) solution (x, y)
  shows a < x ∧ b < y ⟷ pell-valuation (a, b) < pell-valuation (x, y)
proof
  assume a < x ∧ b < y
  thus pell-valuation (a, b) < pell-valuation (x, y)
    unfolding pell-valuation-def prod.case using D-gt-1
    by (intro add-strict-mono mult-strict-right-mono) auto
next
  assume *: pell-valuation (a, b) < pell-valuation (x, y)
  from assms have (a, b) = (x, y) ∨ a < x ∧ b < y ∨ x < a ∧ y < b
    by (rule solutions-linorder-strict)
  thus a < x ∧ b < y
proof (elim disjE)
  assume x < a ∧ y < b
  hence pell-valuation (a, b) > pell-valuation (x, y)
    unfolding pell-valuation-def prod.case using D-gt-1
    by (intro add-strict-mono mult-strict-right-mono) auto
  with * have False by linarith
  thus ?thesis ..
qed (insert *, auto)
qed

```

2.7 The fundamental solution

The *fundamental solution* is the non-trivial solution (x, y) with non-negative x and y for which the Pell valuation $x + y\sqrt{D}$ is minimal, or, equivalently, for which x and y are minimal.

definition *fund-sol* :: $\text{nat} \times \text{nat}$ **where**

fund-sol = (THE $z :: \text{nat} \times \text{nat}$. *is-arg-min* (*pell-valuation* :: $\text{nat} \times \text{nat} \Rightarrow \text{real}$) *nontriv-solution* z)

The well-definedness of this follows from the injectivity of the Pell valuation and the fact that smaller Pell valuation of a solution is smaller than that of another iff the components are both smaller.

theorem *fund-sol-is-arg-min*:

is-arg-min (*pell-valuation* :: $\text{nat} \times \text{nat} \Rightarrow \text{real}$) *nontriv-solution* *fund-sol*

unfolding *fund-sol-def*

proof (rule *theI'*)

show $\exists! z :: \text{nat} \times \text{nat}$. *is-arg-min* (*pell-valuation* :: $\text{nat} \times \text{nat} \Rightarrow \text{real}$) *nontriv-solution* z

proof (rule *ex-ex1I*)

fix $z1\ z2 :: \text{nat} \times \text{nat}$

assume *is-arg-min* (*pell-valuation* :: $\text{nat} \times \text{nat} \Rightarrow \text{real}$) *nontriv-solution* $z1$

is-arg-min (*pell-valuation* :: $\text{nat} \times \text{nat} \Rightarrow \text{real}$) *nontriv-solution* $z2$

hence *pell-valuation* $z1$ = *pell-valuation* $z2$

by (*cases* $z1$, *cases* $z2$, *intro antisym*) (*auto simp: is-arg-min-def not-less*)

thus $z1 = z2$ **by** (*auto split: prod.splits*)

next

define y **where** $y = (\text{LEAST } y. y > 0 \wedge \text{is-square } (1 + D * y^2))$

have $\exists y > 0$. *is-square* $(1 + D * y^2)$

using *pell-solution-exists* **by** (*auto simp: eq-commute[of - Suc -]*)

hence $y: y > 0 \wedge \text{is-square } (1 + D * y^2)$

unfolding *y-def* **by** (rule *LeastI-ex*)

have *y-le*: $y \leq y'$ **if** $y' > 0$ *is-square* $(1 + D * y'^2)$ **for** y'

unfolding *y-def* **using** *that* **by** (*intro Least-le*) *auto*

from y **obtain** x **where** $x: x^2 = 1 + D * y^2$

by (*auto elim: is-nth-powerE*)

with y **have** *nontriv-solution* (x, y)

by (*auto simp: nontriv-solution-def*)

have *is-arg-min* (*pell-valuation* :: $\text{nat} \times \text{nat} \Rightarrow \text{real}$) *nontriv-solution* (x, y)

unfolding *is-arg-min-linorder*

proof *safe*

fix $a\ b :: \text{nat}$

assume $*$: *nontriv-solution* (a, b)

hence $b > 0$ **and** *Suc* $(D * b^2) = a^2$

by (*auto simp: nontriv-solution-def intro!: Nat.gr0I*)

hence *is-square* $(1 + D * b^2)$

by (*auto simp: nontriv-solution-def*)

from $\langle b > 0 \rangle$ **and** *this* **have** $y \leq b$ **by** (rule *y-le*)

```

with ⟨nontriv-solution (x, y)⟩ and * have  $x \leq a$ 
using solutions-linorder-strict[of x y a b] by (auto simp: nontriv-solution-altdef)
with ⟨ $y \leq b$ ⟩ show pell-valuation (int x, int y)  $\leq$  pell-valuation (int a, int b)
unfolding pell-valuation-def prod.case by (intro add-mono mult-right-mono)
auto
qed fact+
thus  $\exists z. \text{is-arg-min } (\text{pell-valuation} :: \text{nat} \times \text{nat} \Rightarrow \text{real}) \text{ nontriv-solution } z ..$ 
qed
qed

```

corollary

```

fund-sol-is-nontriv-solution: nontriv-solution fund-sol
and fund-sol-minimal:
nontriv-solution (a, b)  $\implies$  pell-valuation fund-sol  $\leq$  pell-valuation (int a,
int b)
and fund-sol-minimal':
nontriv-solution (z :: nat  $\times$  nat)  $\implies$  pell-valuation fund-sol  $\leq$  pell-valuation
z
using fund-sol-is-arg-min by (auto simp: is-arg-min-linorder case-prod-unfold)

```

lemma *fund-sol-minimal'':*

```

assumes nontriv-solution z
shows  $\text{fst fund-sol} \leq \text{fst } z \text{ \& } \text{snd fund-sol} \leq \text{snd } z$ 
proof –
have pell-valuation (fst fund-sol, snd fund-sol)  $\leq$  pell-valuation (fst z, snd z)
using fund-sol-minimal'[OF assms] by (simp add: case-prod-unfold)
hence  $\text{fst fund-sol} \leq \text{fst } z \wedge \text{snd fund-sol} \leq \text{snd } z$ 
using assms fund-sol-is-nontriv-solution
by (subst solutions-le-iff-pell-valuation-le) (auto simp: case-prod-unfold)
thus  $\text{fst fund-sol} \leq \text{fst } z \text{ \& } \text{snd fund-sol} \leq \text{snd } z$  by blast+
qed

```

2.8 Group structure on solutions

As was mentioned already, the Pell valuation function provides an injective map from solutions of Pell's equation into the ring $\mathbb{Z}[\sqrt{D}]$. We shall see now that the solutions are actually a subgroup of the multiplicative group of $\mathbb{Z}[\sqrt{D}]$ via the valuation function as a homomorphism:

- The trivial solution $(1, 0)$ has valuation 1 , which is the neutral element of $\mathbb{Z}[\sqrt{D}]^*$
- Multiplication of two solutions $a + b\sqrt{D}$ and $x + y\sqrt{D}$ leads to $\bar{x} + \bar{y}\sqrt{D}$ with $\bar{x} = xa + ybD$ and $\bar{y} = xb + ya$, which is again a solution.
- The conjugate $(x, -y)$ of a solution (x, y) is an inverse element to this multiplication operation, since $(x + y\sqrt{D})(x - y\sqrt{D}) = 1$.

definition *pell-mul* :: ('a :: comm-semiring-1 × 'a) ⇒ ('a × 'a) ⇒ ('a × 'a)
where

pell-mul = (λ(a,b) (x,y). (x * a + y * b * of-nat D, x * b + y * a))

definition *pell-cnj* :: ('a :: comm-ring-1 × 'a) ⇒ 'a × 'a **where**

pell-cnj = (λ(a,b). (a, -b))

lemma *pell-cnj-snd-0* [simp]: *snd* z = 0 ⇒ *pell-cnj* z = z

by (cases z) (simp-all add: *pell-cnj-def*)

lemma *pell-mul-commutes*: *pell-mul* z1 z2 = *pell-mul* z2 z1

by (auto simp: *pell-mul-def algebra-simps case-prod-unfold*)

lemma *pell-mul-assoc*: *pell-mul* z1 (*pell-mul* z2 z3) = *pell-mul* (*pell-mul* z1 z2) z3

by (auto simp: *pell-mul-def algebra-simps case-prod-unfold*)

lemma *pell-mul-trivial-left* [simp]: *pell-mul* (1, 0) z = z

by (auto simp: *pell-mul-def algebra-simps case-prod-unfold*)

lemma *pell-mul-trivial-right* [simp]: *pell-mul* z (1, 0) = z

by (auto simp: *pell-mul-def algebra-simps case-prod-unfold*)

lemma *pell-mul-trivial-left-nat* [simp]: *pell-mul* (Suc 0, 0) z = z

by (auto simp: *pell-mul-def algebra-simps case-prod-unfold*)

lemma *pell-mul-trivial-right-nat* [simp]: *pell-mul* z (Suc 0, 0) = z

by (auto simp: *pell-mul-def algebra-simps case-prod-unfold*)

definition *pell-power* :: ('a :: comm-semiring-1 × 'a) ⇒ nat ⇒ ('a × 'a) **where**

pell-power z n = ((λz'. *pell-mul* z' z) $\overset{\sim}{\sim}$ n) (1, 0)

lemma *pell-power-0* [simp]: *pell-power* z 0 = (1, 0)

by (simp add: *pell-power-def*)

lemma *pell-power-one* [simp]: *pell-power* (1, 0) n = (1, 0)

by (induction n) (auto simp: *pell-power-def*)

lemma *pell-power-one-right* [simp]: *pell-power* z 1 = z

by (simp add: *pell-power-def*)

lemma *pell-power-Suc*: *pell-power* z (Suc n) = *pell-mul* z (*pell-power* z n)

by (simp add: *pell-power-def pell-mul-commutes*)

lemma *pell-power-add*: *pell-power* z (m + n) = *pell-mul* (*pell-power* z m) (*pell-power* z n)

by (induction m arbitrary: z)

(simp-all add: *funpow-add o-def pell-power-Suc pell-mul-assoc*)

lemma *pell-valuation-mult* [simp]:

$pell\text{-}valuation\ (pell\text{-}mul\ z1\ z2) = pell\text{-}valuation\ z1 * pell\text{-}valuation\ z2$
by (*simp add: pell-valuation-def pell-mul-def case-prod-unfold algebra-simps*)

lemma *pell-valuation-mult-nat* [*simp*]:
 $pell\text{-}valuation\ (case\ pell\text{-}mul\ z1\ z2\ of\ (a,\ b) \Rightarrow (int\ a,\ int\ b)) =$
 $pell\text{-}valuation\ z1 * pell\text{-}valuation\ z2$
by (*simp add: pell-valuation-def pell-mul-def case-prod-unfold algebra-simps*)

lemma *pell-valuation-trivial* [*simp*]: $pell\text{-}valuation\ (1,\ 0) = 1$
by (*simp add: pell-valuation-def*)

lemma *pell-valuation-trivial-nat* [*simp*]: $pell\text{-}valuation\ (Suc\ 0,\ 0) = 1$
by (*simp add: pell-valuation-def*)

lemma *pell-valuation-cnj*: $pell\text{-}valuation\ (pell\text{-}cnj\ z) = fst\ z - snd\ z * sqrt\ D$
by (*simp add: pell-valuation-def pell-cnj-def case-prod-unfold*)

lemma *pell-valuation-snd-0* [*simp*]: $pell\text{-}valuation\ (a,\ 0) = of\text{-}int\ a$
by (*simp add: pell-valuation-def*)

lemma *pell-valuation-0-iff* [*simp*]: $pell\text{-}valuation\ z = 0 \longleftrightarrow z = (0,\ 0)$
proof
assume *: $pell\text{-}valuation\ z = 0$
have $snd\ z = 0$
proof (*rule ccontr*)
assume $snd\ z \neq 0$
with * **have** $sqrt\ D = -fst\ z / snd\ z$
by (*simp add: pell-valuation-def case-prod-unfold field-simps*)
also **have** $\dots \in \mathbb{Q}$ **by** *auto*
finally **show** *False* **using** *nonsquare-D irrat-sqrt-nonsquare* **by** *blast*
qed
with * **have** $fst\ z = 0$ **by** (*simp add: pell-valuation-def case-prod-unfold*)
with $\langle snd\ z = 0 \rangle$ **show** $z = (0,\ 0)$ **by** (*cases z*) *auto*
qed (*auto simp: pell-valuation-def*)

lemma *pell-valuation-solution-pos-nat*:
fixes $z :: nat \times nat$
assumes *solution z*
shows $pell\text{-}valuation\ z > 0$
proof –
from *assms* **have** $z \neq (0,\ 0)$ **by** (*intro notI*) *auto*
hence $pell\text{-}valuation\ z \neq 0$ **by** (*auto split: prod.splits*)
moreover **have** $pell\text{-}valuation\ z \geq 0$ **by** (*intro pell-valuation-nonneg*) (*auto split: prod.splits*)
ultimately **show** *?thesis* **by** *linarith*
qed

lemma
assumes *solution z*

shows *pell-mul-cnj-right*: $\text{pell-mul } z (\text{pell-cnj } z) = (1, 0)$
and *pell-mul-cnj-left*: $\text{pell-mul } (\text{pell-cnj } z) z = (1, 0)$
using *assms* **by** (*auto simp: pell-mul-def pell-cnj-def solution-def power2-eq-square*)

lemma *pell-valuation-cnj-solution*:

fixes $z :: \text{nat} \times \text{nat}$

assumes *solution* z

shows $\text{pell-valuation } (\text{pell-cnj } z) = 1 / \text{pell-valuation } z$

proof –

have $\text{pell-valuation } (\text{pell-cnj } z) * \text{pell-valuation } z = \text{pell-valuation } (\text{pell-mul } (\text{pell-cnj } z) z)$

by *simp*

also from *assms* **have** $\text{pell-mul } (\text{pell-cnj } z) z = (1, 0)$

by (*subst pell-mul-cnj-left*) (*auto simp: case-prod-unfold*)

finally show *?thesis* **using** *pell-valuation-solution-pos-nat* [*OF assms*]

by (*auto simp: divide-simps*)

qed

lemma *pell-valuation-power* [*simp*]: $\text{pell-valuation } (\text{pell-power } z n) = \text{pell-valuation } z ^ n$

by (*induction n*) (*simp-all add: pell-power-Suc*)

lemma *pell-valuation-power-nat* [*simp*]:

$\text{pell-valuation } (\text{case } \text{pell-power } z n \text{ of } (a, b) \Rightarrow (\text{int } a, \text{int } b)) = \text{pell-valuation } z ^ n$

by (*induction n*) (*simp-all add: pell-power-Suc*)

lemma *pell-valuation-fund-sol-ge-2*: $\text{pell-valuation } \text{fund-sol} \geq 2$

proof –

obtain $x y$ **where** [*simp*]: $\text{fund-sol} = (x, y)$ **by** (*cases fund-sol*)

from *fund-sol-is-nontriv-solution* **have** $\text{eq}: x^2 = 1 + D * y^2$

by (*auto simp: nontriv-solution-def*)

consider $y > 0 \mid y = 0 \ x \neq 1$

using *fund-sol-is-nontriv-solution* **by** (*force simp: nontriv-solution-def*)

thus *?thesis*

proof *cases*

assume $y > 0$

hence $1 + 1 * 1 \leq 1 + D * y^2$

using *D-pos* **by** (*intro add-mono mult-mono*) *auto*

also from *eq* **have** $\dots = x^2$ **..**

finally have $x^2 > 1^2$ **by** *simp*

hence $x > 1$ **by** (*rule power2-less-imp-less*) *auto*

with $\langle y > 0 \rangle$ **have** $x + y * \text{sqrt } D \geq 2 + 1 * 1$

using *D-pos* **by** (*intro add-mono mult-mono*) *auto*

thus *?thesis* **by** (*simp add: pell-valuation-def*)

next

assume [*simp*]: $y = 0$ **and** $x \neq 1$

with *eq* **have** $x \neq 0$ **by** (*intro notI*) *auto*

with $\langle x \neq 1 \rangle$ **have** $x \geq 2$ **by** *simp*
thus *?thesis* **by** (*auto simp: pell-valuation-def*)
qed
qed

lemma *solution-pell-mul* [*intro*]:
assumes *solution z1 solution z2*
shows *solution (pell-mul z1 z2)*
proof –
obtain *a b* **where** [*simp*]: $z1 = (a, b)$ **by** (*cases z1*)
obtain *c d* **where** [*simp*]: $z2 = (c, d)$ **by** (*cases z2*)
from *assms* **show** *?thesis*
by (*simp add: solution-def pell-mul-def case-prod-unfold power2-eq-square alge-bra-simps*)
qed

lemma *solution-pell-cnj* [*intro*]:
assumes *solution z*
shows *solution (pell-cnj z)*
using *assms* **by** (*auto simp: solution-def pell-cnj-def*)

lemma *solution-pell-power* [*simp, intro*]: *solution z* \implies *solution (pell-power z n)*
by (*induction n*) (*auto simp: pell-power-Suc*)

lemma *pell-mul-eq-trivial-nat-iff*:
 $pell-mul\ z1\ z2 = (Suc\ 0,\ 0) \iff z1 = (Suc\ 0,\ 0) \wedge z2 = (Suc\ 0,\ 0)$
using *D-gt-1* **by** (*cases z1; cases z2*) (*auto simp: pell-mul-def*)

lemma *nontriv-solution-pell-nat-mul1*:
solution (z1 :: nat \times nat) \implies *nontriv-solution z2* \implies *nontriv-solution (pell-mul z1 z2)*
by (*auto simp: nontriv-solution-altdef pell-mul-eq-trivial-nat-iff*)

lemma *nontriv-solution-pell-nat-mul2*:
nontriv-solution (z1 :: nat \times nat) \implies *solution z2* \implies *nontriv-solution (pell-mul z1 z2)*
by (*auto simp: nontriv-solution-altdef pell-mul-eq-trivial-nat-iff*)

lemma *nontriv-solution-power-nat* [*intro*]:
assumes *nontriv-solution (z :: nat \times nat) n > 0*
shows *nontriv-solution (pell-power z n)*
proof –
have *nontriv-solution (pell-power z n) \vee n = 0*
by (*induction n*)
(insert assms(1), auto intro: nontriv-solution-pell-nat-mul1 simp: pell-power-Suc)
with *assms(2)* **show** *?thesis* **by** *auto*
qed

2.9 The different regions of the valuation function

Next, we shall explore what happens to the valuation function for solutions (x, y) for different signs of x and y :

- If $x > 0$ and $y > 0$, we have $x + y\sqrt{D} > 1$.
- If $x > 0$ and $y < 0$, we have $0 < x + y\sqrt{D} < 1$.
- If $x < 0$ and $y > 0$, we have $-1 < x + y\sqrt{D} < 0$.
- If $x < 0$ and $y < 0$, we have $x + y\sqrt{D} < -1$.

In particular, this means that we can deduce the sign of x and y if we know in which of these four regions the valuation lies.

lemma

assumes $x > 0$ $y > 0$ *solution* (x, y)

shows *pell-valuation-pos-pos*: *pell-valuation* $(x, y) > 1$

and *pell-valuation-pos-neg-aux*: *pell-valuation* $(x, -y) \in \{0 < .. < 1\}$

proof –

from *D-gt-1* *assms* **have** $x + y * \text{sqrt } D \geq 1 + 1 * 1$

by (*intro add-mono mult-mono*) *auto*

hence *gt-1*: $x + y * \text{sqrt } D > 1$ **by** *simp*

thus *pell-valuation* $(x, y) > 1$ **by** (*simp add: pell-valuation-def*)

from *assms* **have** $1 = x^2 - D * y^2$ **by** (*simp add: solution-def*)

also **have** *of-int* $\dots = (x - y * \text{sqrt } D) * (x + y * \text{sqrt } D)$

by (*simp add: field-simps power2-eq-square*)

finally **have** *eq*: $(x - y * \text{sqrt } D) = 1 / (x + y * \text{sqrt } D)$

using *gt-1* **by** (*simp add: field-simps*)

note *eq*

also **from** *gt-1* **have** $1 / (x + y * \text{sqrt } D) < 1 / 1$

by (*intro divide-strict-left-mono*) *auto*

finally **have** $x - y * \text{sqrt } D < 1$ **by** *simp*

note *eq*

also **from** *gt-1* **have** $1 / (x + y * \text{sqrt } D) > 0$

by (*intro divide-pos-pos*) *auto*

finally **have** $x - y * \text{sqrt } D > 0$.

with $\langle x - y * \text{sqrt } D < 1 \rangle$ **show** *pell-valuation* $(x, -y) \in \{0 < .. < 1\}$

by (*simp add: pell-valuation-def*)

qed

lemma *pell-valuation-pos-neg*:

assumes $x > 0$ $y < 0$ *solution* (x, y)

shows *pell-valuation* $(x, y) \in \{0 < .. < 1\}$

using *pell-valuation-pos-neg-aux*[*of* $x - y$] *assms* **by** *simp*

lemma *pell-valuation-neg-neg*:
assumes $x < 0$ $y < 0$ *solution* (x, y)
shows *pell-valuation* $(x, y) < -1$
using *pell-valuation-pos-pos*[of $-x -y$] *assms* **by** *simp*

lemma *pell-valuation-neg-pos*:
assumes $x < 0$ $y > 0$ *solution* (x, y)
shows *pell-valuation* $(x, y) \in \{-1 < .. < 0\}$
using *pell-valuation-pos-neg*[of $-x -y$] *assms* **by** *simp*

lemma *pell-valuation-solution-gt1D*:
assumes *solution* z *pell-valuation* $z > 1$
shows $\text{fst } z > 0 \wedge \text{snd } z > 0$
using *pell-valuation-pos-pos*[of $\text{fst } z \text{snd } z$] *pell-valuation-pos-neg*[of $\text{fst } z \text{snd } z$]
pell-valuation-neg-pos[of $\text{fst } z \text{snd } z$] *pell-valuation-neg-neg*[of $\text{fst } z \text{snd } z$]
assms
by (*cases* $\text{fst } z$ 0 :: *int* *rule*: *linorder-cases*;
cases $\text{snd } z$ 0 :: *int* *rule*: *linorder-cases*;
cases z) *auto*

2.10 Generating property of the fundamental solution

We now show that the fundamental solution generates the set of the (non-negative) solutions in the sense that each solution is a power of the fundamental solution. Combined with the symmetry property that (x, y) is a solution iff $(\pm x, \pm y)$ is a solution, this gives us a complete characterisation of all solutions of Pell's equation.

definition *nth-solution* :: $\text{nat} \Rightarrow \text{nat} \times \text{nat}$ **where**
nth-solution $n = \text{pell-power fund-sol } n$

lemma *pell-valuation-nth-solution* [*simp*]:
pell-valuation (*nth-solution* n) = *pell-valuation fund-sol* \wedge^n
by (*simp* *add*: *nth-solution-def*)

theorem *nth-solution-inj*: *inj* *nth-solution*

proof

fix $m\ n$:: *nat*
assume *nth-solution* $m = \text{nth-solution } n$
hence *pell-valuation* (*nth-solution* m) = *pell-valuation* (*nth-solution* n)
by (*simp* *only*:)
also have *pell-valuation* (*nth-solution* m) = *pell-valuation fund-sol* \wedge^m
by *simp*
also have *pell-valuation* (*nth-solution* n) = *pell-valuation fund-sol* \wedge^n
by *simp*
finally show $m = n$
using *pell-valuation-fund-sol-ge-2* **by** (*subst* (*asm*) *power-inject-exp*) *auto*
qed

theorem *nth-solution-sound* [*intro*]: *solution* (*nth-solution* *n*)
using *fund-sol-is-nontriv-solution* **by** (*auto simp: nth-solution-def*)

theorem *nth-solution-sound'* [*intro*]: $n > 0 \implies \text{nontriv-solution}$ (*nth-solution* *n*)
using *fund-sol-is-nontriv-solution* **by** (*auto simp: nth-solution-def*)

theorem *nth-solution-complete*:
fixes $z :: \text{nat} \times \text{nat}$
assumes *solution* z
shows $z \in \text{range } \text{nth-solution}$
proof (*cases* $z = (1, 0)$)
case *True*
hence $z = \text{nth-solution } 0$ **by** (*simp add: nth-solution-def*)
thus *?thesis* **by** *auto*
next
case *False*
with *assms* **have** *nontriv-solution* z **by** (*auto simp: nontriv-solution-altdef*)
show *?thesis*
proof (*rule ccontr*)
assume $\neg ?thesis$
hence $*$: *pell-power fund-sol* $n \neq z$ **for** n **unfolding** *nth-solution-def* **by** *blast*

define u **where** $u = \text{pell-valuation fund-sol}$
define v **where** $v = \text{pell-valuation } z$
define n **where** $n = \text{nat } \lfloor \log u \ v \rfloor$
have *u-ge-2*: $u \geq 2$ **using** *pell-valuation-fund-sol-ge-2* **by** (*auto simp: u-def*)
have *v-pos*: $v > 0$ **unfolding** *v-def* **using** *assms*
by (*intro pell-valuation-solution-pos-nat*) *auto*
have *u-le-v*: $u \leq v$ **unfolding** *u-def v-def* **by** (*rule fund-sol-minimal'*) *fact*

have *u-power-neq-v*: $u^k \neq v$ **for** k
proof
assume $u^k = v$
also **have** $u^k = \text{pell-valuation} (\text{pell-power fund-sol } k)$
by (*simp add: u-def*)
also **have** $\dots = v \iff \text{pell-power fund-sol } k = z$
unfolding *v-def* **by** (*subst pell-valuation-eq-iff*) (*auto split: prod.splits*)
finally **show** *False* **using** $*$ **by** *blast*
qed

from *u-le-v v-pos u-ge-2* **have** *log-ge-1*: $\log u \ v \geq 1$
by (*subst one-le-log-cancel-iff*) *auto*

define z' **where** $z' = \text{pell-mul } z (\text{pell-power} (\text{pell-cn } \text{fund-sol}) \ n)$
define x **and** y **where** $x = \text{nat } \lfloor \text{fst } z' \rfloor$ **and** $y = \text{nat } \lfloor \text{snd } z' \rfloor$
have *solution* z' **using** *assms fund-sol-is-nontriv-solution* **unfolding** *z'-def*
by (*intro solution-pell-mul solution-pell-power solution-pell-cn*) (*auto simp: case-prod-unfold*)

have $u \wedge n < v$
proof –
from *u-ge-2* **have** $u \wedge n = u \text{ powr real } n$ **by** (*subst powr-realpow*) *auto*
also have $\dots \leq u \text{ powr log } u v$ **using** *u-ge-2 log-ge-1*
by (*intro powr-mono*) (*auto simp: n-def*)
also have $\dots = v$
using *u-ge-2 v-pos* **by** (*subst powr-log-cancel*) *auto*
finally have $u \wedge n \leq v$.
with *u-power-neq-v[of n]* **show** *?thesis* **by** *linarith*
qed
moreover have $v < u \wedge \text{Suc } n$
proof –
have $v = u \text{ powr log } u v$
using *u-ge-2 v-pos* **by** (*subst powr-log-cancel*) *auto*
also have $\text{log } u v \leq 1 + \text{real-of-int } \lfloor \text{log } u v \rfloor$ **by** *linarith*
hence $u \text{ powr log } u v \leq u \text{ powr real } (\text{Suc } n)$ **using** *u-ge-2 log-ge-1*
by (*intro powr-mono*) (*auto simp: n-def*)
also have $\dots = u \wedge \text{Suc } n$ **using** *u-ge-2* **by** (*subst powr-realpow*) *auto*
finally have $u \wedge \text{Suc } n \geq v$.
with *u-power-neq-v[of Suc n]* **show** *?thesis* **by** *linarith*
qed
ultimately have $v / u \wedge n \in \{1 <.. < u\}$
using *u-ge-2* **by** (*simp add: field-simps*)
also have $v / u \wedge n = \text{pell-valuation } z'$
using *fund-sol-is-nontriv-solution*
by (*auto simp add: z'-def u-def v-def pell-valuation-cnj-solution field-simps*)
finally have *val: pell-valuation z' ∈ {1 <.. < u}* .

from *val* **and** $\langle \text{solution } z' \rangle$ **have** *nontriv-solution z'*
by (*auto simp: nontriv-solution-altdef*)
from $\langle \text{solution } z' \rangle$ **and** *val* **have** $\text{fst } z' > 0 \wedge \text{snd } z' > 0$
by (*intro pell-valuation-solution-gt1D*) *auto*

hence [*simp*]: $z' = (\text{int } x, \text{int } y)$
by (*auto simp: x-def y-def*)

from $\langle \text{nontriv-solution } z' \rangle$ **have** $\text{pell-valuation } (\text{int } x, \text{int } y) \geq u$
unfolding *u-def* **by** (*intro fund-sol-minimal*) *auto*
with *val* **show** *False* **by** *simp*
qed
qed

corollary *solution-iff-nth-solution*:
fixes $z :: \text{nat} \times \text{nat}$
shows $\text{solution } z \iff z \in \text{range } \text{nth-solution}$
using *nth-solution-sound nth-solution-complete* **by** *blast*

corollary *solution-iff-nth-solution'*:
fixes $z :: \text{int} \times \text{int}$

shows $\text{solution } (a, b) \longleftrightarrow (\text{nat } |a|, \text{nat } |b|) \in \text{range } \text{nth-solution}$
proof –
 have $\text{solution } (a, b) \longleftrightarrow \text{solution } (\text{nat } |a|, \text{nat } |b|)$
 by *simp*
 also have $\dots \longleftrightarrow (\text{nat } |a|, \text{nat } |b|) \in \text{range } \text{nth-solution}$
 by (*rule solution-iff-nth-solution*)
 finally show *?thesis* .
qed

corollary *infinite-solutions*: $\text{infinite } \{z :: \text{nat} \times \text{nat}. \text{solution } z\}$
proof –
 have *infinite (range nth-solution)*
 by (*intro range-inj-infinite nth-solution-inj*)
 also have $\text{range } \text{nth-solution} = \{z :: \text{nat} \times \text{nat}. \text{solution } z\}$
 by (*auto simp: solution-iff-nth-solution*)
 finally show *?thesis* .
qed

corollary *infinite-solutions'*: $\text{infinite } \{z :: \text{int} \times \text{int}. \text{solution } z\}$
proof
 assume *finite {z :: int × int. solution z}*
 hence *finite (map-prod (nat ∘ abs) (nat ∘ abs) ‘ {z :: int × int. solution z})*
 by (*rule finite-imageI*)
 also have $(\text{map-prod } (\text{nat} \circ \text{abs}) (\text{nat} \circ \text{abs}) \text{ ‘ } \{z :: \text{int} \times \text{int}. \text{solution } z\}) =$
 $\{z :: \text{nat} \times \text{nat}. \text{solution } z\}$
 by (*auto simp: map-prod-def image-iff intro!: exI[of - int x for x]*)
 finally show *False using infinite-solutions by contradiction*
qed

lemma *strict-mono-pell-valuation-nth-solution*: $\text{strict-mono } (\text{pell-valuation} \circ \text{nth-solution})$
 using *pell-valuation-fund-sol-ge-2*
 by (*auto simp: strict-mono-def intro!: power-strict-increasing*)

lemma *strict-mono-nth-solution*:
 $\text{strict-mono } (\text{fst} \circ \text{nth-solution}) \text{ strict-mono } (\text{snd} \circ \text{nth-solution})$
proof –
 let *?g = nth-solution*
 have $\text{fst } (?g \ m) < \text{fst } (?g \ n) \wedge \text{snd } (?g \ m) < \text{snd } (?g \ n)$ **if** $m < n$ **for** $m \ n$
 using *pell-valuation-fund-sol-ge-2 that*
 by (*subst solutions-less-iff-pell-valuation-less*) *auto*
 thus $\text{strict-mono } (\text{fst} \circ \text{nth-solution}) \text{ strict-mono } (\text{snd} \circ \text{nth-solution})$
 by (*auto simp: strict-mono-def*)
qed

end

2.11 The case of an “almost square” parameter

If D is equal to $a^2 - 1$ for some $a > 1$, we have a particularly simple case where the fundamental solution is simply $(1, a)$.

context

fixes $a :: nat$
assumes $a: a > 1$

begin

lemma *pell-square-minus1*: $pell (a^2 - Suc\ 0)$

proof

show $\neg is-square (a^2 - Suc\ 0)$

proof

assume $is-square (a^2 - Suc\ 0)$

then obtain k **where** $k^2 = a^2 - 1$ **by** (*auto elim: is-nth-powerE*)

with a **have** $a^2 = Suc (k^2)$ **by** *simp*

hence $a = 1$ **using** *pell-square-solution-nat-iff*[*of 1 a k*] **by** *simp*

with a **show** *False* **by** *simp*

qed

qed

interpretation *pell* $a^2 - Suc\ 0$

by (*rule pell-square-minus1*)

lemma *fund-sol-square-minus1*: $fund-sol = (a, 1)$

proof –

from a **have** $sol: nontriv-solution (a, 1)$

by (*simp add: nontriv-solution-def*)

from sol **have** $snd\ fund-sol \leq 1$

using *fund-sol-minimal'*[*of (a, 1)*] **by** *auto*

with *solutions-linorder-strict*[*of a 1 fst fund-sol snd fund-sol*]

fund-sol-is-nontriv-solution sol

show $fund-sol = (a, 1)$

by (*cases fund-sol*) (*auto simp: nontriv-solution-altdef*)

qed

end

2.12 Alternative presentation of the main results

theorem *pell-solutions*:

fixes $D :: nat$

assumes $\nexists k. D = k^2$

obtains $x_0\ y_0 :: nat$

where $\forall (x::int) (y::int).$

$x^2 - D * y^2 = 1 \iff$

$(\exists n::nat. nat\ |x| + sqrt\ D * nat\ |y| = (x_0 + sqrt\ D * y_0) \wedge n)$

proof –

from *assms* **interpret** *pell*

```

    by unfold-locales (auto simp: is-nth-power-def)
  show ?thesis
proof (rule that[of fst fund-sol snd fund-sol], intro allI, goal-cases)
  case (1 x y)
  have  $(x^2 - \text{int } D * y^2 = 1) \longleftrightarrow \text{solution } (x, y)$ 
    by (auto simp: solution-def)
  also have  $\dots \longleftrightarrow (\exists n. (\text{nat } |x|, \text{nat } |y|) = \text{nth-solution } n)$ 
    by (subst solution-iff-nth-solution) blast
  also have  $(\lambda n. (\text{nat } |x|, \text{nat } |y|) = \text{nth-solution } n) =$ 
     $(\lambda n. \text{pell-valuation } (\text{nat } |x|, \text{nat } |y|) = \text{pell-valuation } (\text{nth-solution } n))$ 
    by (subst pell-valuation-eq-iff) (auto simp add: case-prod-unfold prod-eq-iff)
  fun-eq-iff
  also have  $\dots = (\lambda n. \text{nat } |x| + \text{sqrt } D * \text{nat } |y| = (\text{fst fund-sol} + \text{sqrt } D * \text{snd}$ 
     $\text{fund-sol}) \wedge n)$ 
    by (subst pell-valuation-nth-solution)
    (simp add: pell-valuation-def case-prod-unfold mult-ac)
  finally show ?case .
qed
qed

```

```

corollary pell-solutions-infinite:
  fixes D :: nat
  assumes  $\nexists k. D = k^2$ 
  shows infinite  $\{(x :: \text{int}, y :: \text{int}). x^2 - D * y^2 = 1\}$ 
proof -
  from assms interpret pell
    by unfold-locales (auto simp: is-nth-power-def)
  have  $\{(x :: \text{int}, y :: \text{int}). x^2 - D * y^2 = 1\} = \{z. \text{solution } z\}$ 
    by (auto simp: solution-def)
  also have infinite  $\dots$  by (rule infinite-solutions)
  finally show ?thesis .
qed
end

```

2.13 Executable code

```

theory Pell-Algorithm
imports
  Pell
  Efficient-Discrete-Sqrt
  HOL-Library.Discrete
  HOL-Library.While-Combinator
  HOL-Library.Stream
begin

```

2.13.1 Efficient computation of powers by squaring

The following is a tail-recursive implementation of exponentiation by squaring. It works for any binary operation f that fulfils $f x (f x z) = f (f x x) z$, i. e. some weak form of associativity.

context

fixes $f :: 'a \Rightarrow 'a \Rightarrow 'a$

begin

function $efficient-power :: 'a \Rightarrow 'a \Rightarrow nat \Rightarrow 'a$ **where**

$efficient-power\ y\ x\ 0 = y$

| $efficient-power\ y\ x\ (Suc\ 0) = f\ x\ y$

| $n \neq 0 \implies even\ n \implies efficient-power\ y\ x\ n = efficient-power\ y\ (f\ x\ x)\ (n\ div\ 2)$

| $n \neq 1 \implies odd\ n \implies efficient-power\ y\ x\ n = efficient-power\ (f\ x\ y)\ (f\ x\ x)\ (n\ div\ 2)$

by $force+$

termination by $(relation\ measure\ (snd\ \circ\ snd))\ (auto\ elim:\ oddE)$

lemma $efficient-power-code$ [code]:

$efficient-power\ y\ x\ n =$

(if $n = 0$ then y

else if $n = 1$ then $f\ x\ y$

else if even n then $efficient-power\ y\ (f\ x\ x)\ (n\ div\ 2)$

else $efficient-power\ (f\ x\ y)\ (f\ x\ x)\ (n\ div\ 2)$)

by $(induction\ y\ x\ n\ rule:\ efficient-power.induct)\ auto$

lemma $efficient-power-correct$:

assumes $\bigwedge x\ z. f\ x\ (f\ x\ z) = f\ (f\ x\ x)\ z$

shows $efficient-power\ y\ x\ n = (f\ x\ \overset{\sim}{\sim} n)\ y$

proof –

have [simp]: $f\ \overset{\sim}{\sim} 2 = (\lambda x. f\ (f\ x))$ **for** $f :: 'a \Rightarrow 'a$

by $(simp\ add:\ eval-nat-numeral\ o-def)$

show ?thesis

by $(induction\ y\ x\ n\ rule:\ efficient-power.induct)$

$(auto\ elim!\ evenE\ oddE\ simp:\ funpow-mult\ [symmetric]\ funpow-Suc-right$

$assms$

$simp\ del:\ funpow.simps(2))$

qed

end

2.13.2 Multiplication and powers of solutions

We define versions of Pell solution multiplication and exponentiation specialised to natural numbers, both for efficiency reasons and to circumvent the problem of generating code for definitions made inside locales.

fun $pell-mul-nat :: nat \Rightarrow nat \times nat \Rightarrow -$ **where**

$pell-mul-nat\ D\ (a,\ b)\ (x,\ y) = (a * x + D * b * y,\ a * y + b * x)$

lemma (in *pell*) *pell-mul-nat-correct* [*simp*]: *pell-mul-nat* $D = \text{pell.pell-mul } D$
by (*auto simp add: pell-mul-def fun-eq-iff*)

definition *efficient-pell-power* :: $\text{nat} \Rightarrow \text{nat} \times \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \times \text{nat}$ **where**
efficient-pell-power $D z n = \text{efficient-power } (\text{pell-mul-nat } D) (1, 0) z n$

lemma *efficient-pell-power-correct* [*simp*]:
efficient-pell-power $D z n = (\text{pell-mul-nat } D z \hat{\sim} n) (1, 0)$
unfolding *efficient-pell-power-def*
by (*intro efficient-power-correct*) (*auto simp: algebra-simps*)

2.13.3 Finding the fundamental solution

In the following, we set up a very simple algorithm for computing the fundamental solution (x, y) . We try increasing values for y until $1 + Dy^2$ is a perfect square, which we check using an efficient square-detection algorithm. This is efficient enough to work on some interesting small examples.

Much better algorithms (typically based on the continued fraction expansion of \sqrt{D}) are available, but they are also considerably more complicated.

lemma *Discrete-sqrt-square-is-square*:
assumes *is-square* n
shows *Discrete.sqrt* $n \wedge 2 = n$
using *assms* **unfolding** *is-nth-power-def* **by** *force*

definition *find-fund-sol-step* :: $\text{nat} \Rightarrow \text{nat} \times \text{nat} + \text{nat} \times \text{nat} \Rightarrow -$ **where**
find-fund-sol-step $D = (\lambda \text{Inl } (y, y') \Rightarrow$
(case get-nat-sqrt y' *of*
Some $x \Rightarrow \text{Inr } (x, y)$
 $|\ \text{None} \Rightarrow \text{Inl } (y + 1, y' + D * (2 * y + 1))$ *))*

definition *find-fund-sol* **where**
find-fund-sol $D =$
(if square-test D *then*
 $(0, 0)$
else
 $\text{sum.projr } (\text{while sum.isl } (\text{find-fund-sol-step } D) (\text{Inl } (1, 1 + D)))$ *)*

lemma *fund-sol-code*:
assumes $\neg \text{is-square } (D :: \text{nat})$
shows *pell.fund-sol* $D = \text{sum.projr } (\text{while isl } (\text{find-fund-sol-step } D) (\text{Inl } (\text{Suc } 0, \text{Suc } D)))$
proof –
from *assms* **interpret** *pell* D **by** *unfold-locales*
note [*simp*] = *find-fund-sol-step-def*
define f **where** $f = \text{find-fund-sol-step } D$
define $P :: \text{nat} \Rightarrow \text{bool}$ **where** $P = (\lambda y. y > 0 \wedge \text{is-square } (y^2 * D + 1))$
define $Q :: \text{nat} \times \text{nat} \Rightarrow \text{bool}$ **where**

$Q = (\lambda(x,y). P y \wedge (\forall y' \in \{0 < .. < y\}. \neg P y') \wedge x = \text{Discrete.sqrt } (y^2 * D + 1))$

define $R :: \text{nat} \times \text{nat} + \text{nat} \times \text{nat} \Rightarrow \text{bool}$
where $R = (\lambda s. \text{case } s \text{ of}$
 $\quad \text{Inl } (m, m') \Rightarrow m > 0 \wedge (m' = m^2 * D + 1) \wedge (\forall y \in \{0 < .. < m\}. \neg \text{is-square } (y^2 * D + 1))$
 $\quad | \text{Inr } x \Rightarrow Q x)$

define $\text{rel} :: ((\text{nat} \times \text{nat} + \text{nat} \times \text{nat}) \times (\text{nat} \times \text{nat} + \text{nat} \times \text{nat})) \text{ set}$
where $\text{rel} = \{(A,B). (\text{case } (A, B) \text{ of}$
 $\quad (\text{Inl } (m, -), \text{Inl } (m', -)) \Rightarrow m' > 0 \wedge m > m' \wedge m \leq \text{snd } \text{fund-sol}$
 $\quad | (\text{Inr } -, \text{Inl } (m', -)) \Rightarrow m' \leq \text{snd } \text{fund-sol}$
 $\quad | - \Rightarrow \text{False}) \wedge A = f B\}$

obtain $x y$ **where** $xy: \text{sum.projr } (\text{while } \text{isl } f (\text{Inl } (\text{Suc } 0, \text{Suc } D))) = (x, y)$
by $(\text{cases } \text{sum.projr } (\text{while } \text{isl } f (\text{Inl } (\text{Suc } 0, \text{Suc } D))))$

have $\text{neq-fund-solI}: y \neq \text{snd } \text{fund-sol}$ **if** $\neg \text{is-square } (\text{Suc } (y^2 * D))$ **for** y
proof
assume $y = \text{snd } \text{fund-sol}$
with $\text{fund-sol-is-nontriv-solution}$ **have** $\text{Suc } (y^2 * D) = \text{fst } \text{fund-sol} \wedge 2$
by $(\text{simp add: nontriv-solution-def case-prod-unfold})$
hence $\text{is-square } (\text{Suc } (y^2 * D))$ **by** simp
with that show False **by** contradiction
qed

have $\text{case-sum } (\lambda -. \text{False}) Q (\text{while } \text{sum.isl } f (\text{Inl } (m, m^2 * D + 1)))$
if $\forall y \in \{0 < .. < m\}. \neg \text{is-square } (y^2 * D + 1) \wedge m > 0$ **for** m
proof $(\text{rule } \text{while-rule}[\text{where } b = \text{sum.isl}])$
show $R (\text{Inl } (m, m^2 * D + 1))$
using that by $(\text{auto simp: } R\text{-def})$

next
fix s **assume** $R s \text{isl } s$
thus $R (f s)$
by $(\text{auto simp: not-less-less-Suc-eq } Q\text{-def } P\text{-def } R\text{-def } f\text{-def } \text{get-nat-sqrt-def } \text{power2-eq-square algebra-simps split: sum.splits prod.splits})$

next
fix s **assume** $R s \neg \text{isl } s$
thus $\text{case } s \text{ of } \text{Inl } - \Rightarrow \text{False} \mid \text{Inr } x \Rightarrow Q x$
by $(\text{auto simp: } R\text{-def split: sum.splits})$

next
fix s **assume** $s: R s \text{isl } s$
show $(f s, s) \in \text{rel}$
proof $(\text{cases } s)$
case $[\text{simp}]: (\text{Inl } s')$
obtain $a b$ **where** $[\text{simp}]: s' = (a, b)$ **by** $(\text{cases } s')$
from s **have** $*$: $a > 0 \wedge b = \text{Suc } (a^2 * D) \wedge y. y \in \{0 < .. < a\} \implies \neg \text{is-square } (\text{Suc } (y^2 * D))$
by $(\text{auto simp: } R\text{-def})$

```

have a < snd fund-sol if **: ¬ is-square (Suc (a2 * D))
proof -
  from neq-fund-solI have y' ≠ snd fund-sol if y' ∈ {0 <.. $Suc$  a} for y'
  using * ** that by (cases y' = a) auto
  moreover have snd fund-sol ≠ 0 using fund-sol-is-nontriv-solution
  by (intro notI, cases fund-sol) (auto simp: nontriv-solution-altdef)
  ultimately have ∀ y' ≤ a. y' ≠ snd fund-sol by (auto simp: less-Suc-eq-le)
  thus snd fund-sol > a by (cases a < snd fund-sol) (auto simp: not-less)
qed
moreover have a ≤ snd fund-sol
proof -
  have ∀ y' ∈ {0 <.. $a$ }. y' ≠ snd fund-sol using neq-fund-solI *
  by (auto simp: less-Suc-eq-le)
  moreover have snd fund-sol ≠ 0 using fund-sol-is-nontriv-solution
  by (intro notI, cases fund-sol) (auto simp: nontriv-solution-altdef)
  ultimately have ∀ y' < a. y' ≠ snd fund-sol by (auto simp: less-Suc-eq-le)
  thus snd fund-sol ≥ a by (cases a ≤ snd fund-sol) (auto simp: not-less)
qed
ultimately show ?thesis using *
  by (auto simp: f-def get-nat-sqrt-def rel-def)
qed (insert s, auto)
next
define rel'
  where rel' = {(y, x). (case x of Inl (m, -) ⇒ m ≤ snd fund-sol | Inr - ⇒
False) ∧ y = f x}
  have wf rel' unfolding rel'-def
  by (rule wf-if-measure[where f = λz. case z of Inl (m, -) ⇒ Suc (snd fund-sol)
- m | - ⇒ 0])
  (auto split: prod.splits sum.splits simp: f-def get-nat-sqrt-def)
  moreover have rel ⊆ rel'
  proof safe
    fix w z assume (w, z) ∈ rel
    thus (w, z) ∈ rel' by (cases w; cases z) (auto simp: rel-def rel'-def)
  qed
  ultimately show wf rel by (rule wf-subset)
qed
from this[of 1] and xy have *: Q (x, y)
  by (auto split: sum.splits)

from * have is-square (Suc (y2 * D)) by (simp add: Q-def P-def)
with * have x2 = Suc (y2 * D) y > 0
  by (auto simp: Q-def P-def Discrete-sqrt-square-is-square)
hence nontriv-solution (x, y)
  by (auto simp: nontriv-solution-def)
from this have snd fund-sol ≤ snd (x, y)
  by (rule fund-sol-minimal'')
moreover have snd fund-sol ≥ y
proof -
  from * have (∀ y' ∈ {0 <.. $y$ }. ¬ is-square (Suc (y'2 * D)))

```

```

    by (simp add: Q-def P-def)
  with neq-fund-solI have (∀ y' ∈ {0 <..

```

```

lemma find-fund-sol-correct: find-fund-sol D = (if is-square D then (0, 0) else
pell.fund-sol D)
  by (simp add: find-fund-sol-def fund-sol-code square-test-correct)

```

2.13.4 The infinite list of all solutions

```

definition pell-solutions :: nat ⇒ (nat × nat) stream where
  pell-solutions D = (let z = find-fund-sol D in siterate (pell-mul-nat D z) (1, 0))

```

```

lemma (in pell) snth (pell-solutions D) n = nth-solution n
  by (simp add: pell-solutions-def Let-def find-fund-sol-correct nonsquare-D nth-solution-def
pell-power-def pell-mul-commutes[of - fund-sol])

```

2.13.5 Computing the n -th solution

```

definition find-nth-solution :: nat ⇒ nat ⇒ nat × nat where
  find-nth-solution D n =
    (if is-square D then (0, 0) else
    let z = sum.projr (while isl (find-fund-sol-step D) (Inl (Suc 0, Suc D)))
    in efficient-pell-power D z n)

```

```

lemma (in pell) find-nth-solution-correct: find-nth-solution D n = nth-solution n
  by (simp add: find-nth-solution-def nonsquare-D nth-solution-def fund-sol-code
pell-power-def pell-mul-commutes[of - projr -])

```

end

2.13.6 Tests

```

theory Pell-Algorithm-Test
imports
  Pell-Algorithm

```



```
HOL-Library.Code-Target-Numeral
HOL-Library.Code-Lazy
begin

code-lazy-type stream

value find-fund-sol 73
value find-fund-sol 106

value stake 100 (pell-solutions 73)
value snth (pell-solutions 73) 600

value find-nth-solution 73 600
value find-nth-solution 106 10

end
```

References

- [1] Pell's equation, handout for MATHS 714. Lecture notes, University of Auckland, 2008.
- [2] H. Cohen. *A Course in Computational Algebraic Number Theory*. Springer, 2010.
- [3] M. Jacobson and H. Williams. *Solving the Pell Equation*. CMS Books in Mathematics. Springer New York, 2008.