

Ordinary Differential Equations

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Abstract

Session `Ordinary-Differential-Equations` formalizes ordinary differential equations (ODEs) and initial value problems. This work comprises proofs for local and global existence of unique solutions (Picard-Lindelöf theorem). Moreover, it contains a formalization of the (continuous or even differentiable) dependency of the flow on initial conditions as the *flow* of ODEs.

Not in the generated document are the following sessions:

- `HOL-ODE-Numerics`: Rigorous numerical algorithms for computing enclosures of solutions based on Runge-Kutta methods and affine arithmetic. Reachability analysis with splitting and reduction at hyperplanes.
- `HOL-ODE-Examples`: Applications of the numerical algorithms to concrete systems of ODEs (e.g., van der Pol and Lorenz attractor).

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1 Auxiliary Lemmas

theory *ODE-Auxiliarities*

imports

HOL-Analysis.Analysis

HOL-Library.Float

List-Index.List-Index

Affine-Arithmetic.Affine-Arithmetic-Auxiliarities

Affine-Arithmetic.Executable-Euclidean-Space

begin

instantiation *prod* :: (*zero-neq-one*, *zero-neq-one*) *zero-neq-one*

begin

definition *1* = (*1*, *1*)

instance *<proof>*

end

1.1 there is no inner product for type $'a \Rightarrow_L 'b$

lemma (*in real-inner*) *parallelogram-law*: $(\text{norm } (x + y))^2 + (\text{norm } (x - y))^2 = 2 * (\text{norm } x)^2 + 2 * (\text{norm } y)^2$

<proof>

locale *no-real-inner*

begin

lift-definition *fstzero*::(*real*real*) \Rightarrow_L (*real*real*) **is** $\lambda(x, y). (x, 0)$

<proof>

lemma [*simp*]: *fstzero* (*a*, *b*) = (*a*, *0*)

<proof>

lift-definition *zerosnd*::(*real*real*) \Rightarrow_L (*real*real*) **is** $\lambda(x, y). (0, y)$

<proof>

lemma [*simp*]: *zerosnd* (*a*, *b*) = (*0*, *b*)

<proof>

lemma *fstzero-add-zerosnd*: *fstzero* + *zerosnd* = *id-blifun*

<proof>

lemma *norm-fstzero-zerosnd*: *norm* *fstzero* = *1* *norm* *zerosnd* = *1* *norm* (*fstzero* - *zerosnd*) = *1*

<proof>

compare with $(\text{norm } (?x + ?y))^2 + (\text{norm } (?x - ?y))^2 = 2 * (\text{norm } ?x)^2 + 2 * (\text{norm } ?y)^2$

lemma $(\text{norm } (\text{fstzero} + \text{zerosnd}))^2 + (\text{norm } (\text{fstzero} - \text{zerosnd}))^2 \neq$
 $2 * (\text{norm } \text{fstzero})^2 + 2 * (\text{norm } \text{zerosnd})^2$
 ⟨proof⟩

end

1.2 Topology

1.3 Vector Spaces

lemma *ex-norm-eq-1*: $\exists x. \text{norm } (x::'a::\{\text{real-normed-vector}, \text{perfect-space}\}) = 1$
 ⟨proof⟩

1.4 Reals

1.5 Balls

sometimes $(?y \in \text{ball } ?x ?e) = (\text{dist } ?x ?y < ?e)$ etc. are not good [*simp*] rules (although they are often useful): not sure that inequalities are “simpler” than set membership (distorts automatic reasoning when only sets are involved)

lemmas [*simp del*] = *mem-ball mem-cball mem-sphere mem-ball-0 mem-cball-0*

1.6 Boundedness

lemma *bounded-subset-cboxE*:
assumes $\bigwedge i. i \in \text{Basis} \implies \text{bounded } ((\lambda x. x \cdot i) ' X)$
obtains *a b where* $X \subseteq \text{cbox } a b$
 ⟨proof⟩

lemma
bounded-euclideanI:
assumes $\bigwedge i. i \in \text{Basis} \implies \text{bounded } ((\lambda x. x \cdot i) ' X)$
shows *bounded* X
 ⟨proof⟩

1.7 Intervals

notation *closed-segment* $\langle (1\{\text{---}\}) \rangle$
notation *open-segment* $\langle (1\{\text{<--<-}\}) \rangle$

lemma *min-zero-mult-nonneg-le*: $0 \leq h' \implies h' \leq h \implies \text{min } 0 (h * k::\text{real}) \leq h' * k$
 ⟨proof⟩

lemma *max-zero-mult-nonneg-le*: $0 \leq h' \implies h' \leq h \implies h' * k \leq \text{max } 0 (h * k::\text{real})$
 ⟨proof⟩

lemmas *closed-segment-eq-real-ivl = closed-segment-eq-real-ivl*

lemma *bdd-above-is-interval**I*: *bdd-above I if is-interval I a ≤ b a ∈ I b ∉ I for I::real set*
 ⟨proof⟩

lemma *bdd-below-is-interval**I*: *bdd-below I if is-interval I a ≤ b a ∉ I b ∈ I for I::real set*
 ⟨proof⟩

1.8 Extended Real Intervals

1.9 Euclidean Components

1.10 Operator Norm

1.11 Limits

lemma *eventually-open-cball*:
 assumes *open X*
 assumes *x ∈ X*
 shows *eventually (λe. cball x e ⊆ X) (at-right 0)*
 ⟨proof⟩

1.12 Continuity

1.13 Derivatives

lemma
if-eventually-has-derivative:
 assumes *(f has-derivative F') (at x within S)*
 assumes $\forall_F x$ *in at x within S. P x P x x ∈ S*
 shows *((λx. if P x then f x else g x) has-derivative F') (at x within S)*
 ⟨proof⟩

lemma *norm-le-in-cube**I*: *norm x ≤ norm y*
 if $\bigwedge i. i \in \text{Basis} \implies \text{abs } (x \cdot i) \leq \text{abs } (y \cdot i)$ **for** *x y*
 ⟨proof⟩

lemma *has-derivative-partials-euclidean-convex**I*:
 fixes *f::'a::euclidean-space ⇒ 'b::real-normed-vector*
 assumes *f'*: $\bigwedge i x xi. i \in \text{Basis} \implies (\forall j \in \text{Basis}. x \cdot j \in X j) \implies xi = x \cdot i \implies ((\lambda p. f (x + (p - x \cdot i) *_R i)) \text{ has-vector-derivative } f' i x)$ *(at xi within X i)*
 assumes *df-cont*: $\bigwedge i. i \in \text{Basis} \implies (f' i \longrightarrow (f' i x))$ *(at x within {x. ∀ j ∈ Basis. x · j ∈ X j})*
 assumes $\bigwedge i. i \in \text{Basis} \implies x \cdot i \in X i$
 assumes $\bigwedge i. i \in \text{Basis} \implies \text{convex } (X i)$
 shows *(f has-derivative (λh. ∑ j ∈ Basis. (h · j) *_R f' j x)) (at x within {x. ∀ j ∈ Basis. x · j ∈ X j})*
 (is - *(at x within ?S)*)
 ⟨proof⟩

lemma*frechet-derivative-equals-partial-derivative:***fixes** $f :: 'a :: euclidean-space \Rightarrow 'a$ **assumes** $Df: \bigwedge x. (f \text{ has-derivative } Df \ x) \text{ (at } x)$ **assumes** $f': ((\lambda p. f \ (x + (p - x \cdot i) *_{\mathbb{R}} i) \cdot b) \text{ has-real-derivative } f' \ x \ i \ b) \text{ (at } (x \cdot i))$ **shows** $Df \ x \ i \cdot b = f' \ x \ i \ b$ *<proof>*

1.14 Integration

lemmas *content-real[simp]***lemmas** *integrable-continuous[intro, simp]***and** *integrable-continuous-real[intro, simp]***lemma** *integral-eucl-le:***fixes** $f \ g :: 'a :: euclidean-space \Rightarrow 'b :: ordered-euclidean-space$ **assumes** $f \text{ integrable-on } s$ **and** $g \text{ integrable-on } s$ **and** $\bigwedge x. x \in s \implies f \ x \leq g \ x$ **shows** $\text{integral } s \ f \leq \text{integral } s \ g$ *<proof>***lemma***integral-ivl-bound:***fixes** $l \ u :: 'a :: ordered-euclidean-space$ **assumes** $\bigwedge x \ h'. h' \in \{t0 \ .. \ h\} \implies x \in \{t0 \ .. \ h\} \implies (h' - t0) *_{\mathbb{R}} f \ x \in \{l \ .. \ u\}$ **assumes** $t0 \leq h$ **assumes** $f\text{-int}: f \text{ integrable-on } \{t0 \ .. \ h\}$ **shows** $\text{integral } \{t0 \ .. \ h\} \ f \in \{l \ .. \ u\}$ *<proof>***lemma***add-integral-ivl-bound:***fixes** $l \ u :: 'a :: ordered-euclidean-space$ **assumes** $\bigwedge x \ h'. h' \in \{t0 \ .. \ h\} \implies x \in \{t0 \ .. \ h\} \implies (h' - t0) *_{\mathbb{R}} f \ x \in \{l - x0 \ .. \ u - x0\}$ **assumes** $t0 \leq h$ **assumes** $f\text{-int}: f \text{ integrable-on } \{t0 \ .. \ h\}$ **shows** $x0 + \text{integral } \{t0 \ .. \ h\} \ f \in \{l \ .. \ u\}$ *<proof>*

1.15 conditionally complete lattice

1.16 Lists

lemma*Ball-set-Cons[simp]:* $(\forall a \in \text{set-Cons } x \ y. P \ a) \iff (\forall a \in x. \forall b \in y. P \ (a \# b))$

<proof>

lemma *set-cons-eq-empty*[iff]: *set-Cons a b = {}* \longleftrightarrow *a = {}* \vee *b = {}*
<proof>

lemma *listset-eq-empty-iff*[iff]: *listset XS = {}* \longleftrightarrow *{}* \in *set XS*
<proof>

lemma *sing-in-sings*[simp]: *[x] \in ($\lambda x. [x]$) 'xd* \longleftrightarrow *x \in xd*
<proof>

lemma *those-eq-None-set-iff*: *those xs = None* \longleftrightarrow *None \in set xs*
<proof>

lemma *those-eq-Some-lengthD*: *those xs = Some ys* \implies *length xs = length ys*
<proof>

lemma *those-eq-Some-map-Some-iff*: *those xs = Some ys* \longleftrightarrow (*xs = map Some ys*) (*is ?l* \longleftrightarrow *?r*)
<proof>

1.17 Set(sum)

1.18 Max

1.19 Uniform Limit

1.20 Bounded Linear Functions

lift-definition *comp3*::— TODO: name?
(*'c::real-normed-vector \Rightarrow_L 'd::real-normed-vector*) \Rightarrow (*'b::real-normed-vector \Rightarrow_L 'c*) \Rightarrow_L (*'b \Rightarrow_L 'd*) **is**
 $\lambda(cd::('c \Rightarrow_L 'd)) (bc::'b \Rightarrow_L 'c). (cd \circ_L bc)$
<proof>

lemma *blinfun-apply-comp3*[simp]: *blinfun-apply (comp3 a) b = (a \circ_L b)*
<proof>

lemma *bounded-linear-comp3*[bounded-linear]: *bounded-linear comp3*
<proof>

lift-definition *comp12*::— TODO: name?
(*'a::real-normed-vector \Rightarrow_L 'c::real-normed-vector*) \Rightarrow (*'b::real-normed-vector \Rightarrow_L 'c*) \Rightarrow (*'a \times 'b*) \Rightarrow_L *'c*
is $\lambda f g (a, b). f a + g b$
<proof>

lemma *blinfun-apply-comp12*[simp]: *blinfun-apply (comp12 f g) b = f (fst b) + g (snd b)*
<proof>

1.21 Order Transitivity Attributes

$\langle ML \rangle$

1.22 point reflection

definition $\text{preflect}::'a::\text{real-vector} \Rightarrow 'a \Rightarrow 'a$ **where** $\text{preflect} \equiv \lambda t0 t. 2 *_{\mathbb{R}} t0 - t$

lemma $\text{preflect-preflect}[simp]: \text{preflect } t0 (\text{preflect } t0 t) = t$
 $\langle proof \rangle$

lemma $\text{preflect-preflect-image}[simp]: \text{preflect } t0 ` \text{preflect } t0 ` S = S$
 $\langle proof \rangle$

lemma $\text{is-interval-preflect}[simp]: \text{is-interval } (\text{preflect } t0 ` S) \longleftrightarrow \text{is-interval } S$
 $\langle proof \rangle$

lemma $\text{iv-in-preflect-image}[intro, simp]: t0 \in T \Longrightarrow t0 \in \text{preflect } t0 ` T$
 $\langle proof \rangle$

lemma $\text{preflect-tendsto}[tendsto-intros]:$
fixes $l::'a::\text{real-normed-vector}$
shows $(g \longrightarrow l) F \Longrightarrow (h \longrightarrow m) F \Longrightarrow ((\lambda x. \text{preflect } (g x) (h x)) \longrightarrow \text{preflect } l m) F$
 $\langle proof \rangle$

lemma $\text{continuous-preflect}[continuous-intros]:$
fixes $a::'a::\text{real-normed-vector}$
shows $\text{continuous } (\text{at } a \text{ within } A) (\text{preflect } t0)$
 $\langle proof \rangle$

lemma
fixes $t0::'a::\text{ordered-real-vector}$
shows $\text{preflect-le}[simp]: t0 \leq \text{preflect } t0 b \longleftrightarrow b \leq t0$
and $\text{le-preflect}[simp]: \text{preflect } t0 b \leq t0 \longleftrightarrow t0 \leq b$
and $\text{antimono-preflect}: \text{antimono } (\text{preflect } t0)$
and $\text{preflect-le-preflect}[simp]: \text{preflect } t0 a \leq \text{preflect } t0 b \longleftrightarrow b \leq a$
and $\text{preflect-eq-cancel}[simp]: \text{preflect } t0 a = \text{preflect } t0 b \longleftrightarrow a = b$
 $\langle proof \rangle$

lemma $\text{preflect-eq-point-iff}[simp]: t0 = \text{preflect } t0 s \longleftrightarrow t0 = s \text{ preflect } t0 s = t0$
 $\longleftrightarrow t0 = s$
 $\langle proof \rangle$

lemma $\text{preflect-minus-self}[simp]: \text{preflect } t0 s - t0 = t0 - s$
 $\langle proof \rangle$

end

theory $MVT-Ex$

imports

HOL-Analysis.Analysis
HOL-Decision-Procs.Approximation
../ODE-Auxiliarities

begin

1.23 (Counter)Example of Mean Value Theorem in Euclidean Space

There is no exact analogon of the mean value theorem in the multivariate case!

lemma *MVT-wrong: assumes*

$\bigwedge J a u (f::real*real \Rightarrow real*real).$
 $(\bigwedge x. FDERIV f x :> J x) \Longrightarrow$
 $(\exists t \in \{0 <..< 1\}. f (a + u) - f a = J (a + t *R u) u)$

shows *False*

<proof>

lemma *MVT-corrected:*

fixes $f::'a::ordered-euclidean-space \Rightarrow 'b::euclidean-space$
assumes $fderiv: \bigwedge x. x \in D \Longrightarrow (f \text{ has-derivative } J x) \text{ (at } x \text{ within } D)$
assumes $line-in: \bigwedge x. \llbracket 0 \leq x; x \leq 1 \rrbracket \Longrightarrow a + x *R u \in D$
shows $(\exists t \in Basis \rightarrow \{0 <..< 1\}. (f (a + u) - f a) = (\sum i \in Basis. (J (a + t i *R u) u \cdot i) *R i))$

<proof>

lemma *MVT-ivl:*

fixes $f::'a::ordered-euclidean-space \Rightarrow 'b::ordered-euclidean-space$
assumes $fderiv: \bigwedge x. x \in D \Longrightarrow (f \text{ has-derivative } J x) \text{ (at } x \text{ within } D)$
assumes $J-ivl: \bigwedge x. x \in D \Longrightarrow J x u \in \{J0 .. J1\}$
assumes $line-in: \bigwedge x. x \in \{0..1\} \Longrightarrow a + x *R u \in D$
shows $f (a + u) - f a \in \{J0..J1\}$

<proof>

lemma *MVT:*

shows
 $\bigwedge J J0 J1 a u (f::real*real \Rightarrow real*real).$
 $(\bigwedge x. FDERIV f x :> J x) \Longrightarrow$
 $(\bigwedge x. J x u \in \{J0 .. J1\}) \Longrightarrow$
 $f (a + u) - f a \in \{J0 .. J1\}$

<proof>

lemma *MVT-ivl':*

fixes $f::'a::ordered-euclidean-space \Rightarrow 'b::ordered-euclidean-space$
assumes $fderiv: (\bigwedge x. x \in D \Longrightarrow (f \text{ has-derivative } J x) \text{ (at } x \text{ within } D))$
assumes $J-ivl: \bigwedge x. x \in D \Longrightarrow J x (a - b) \in \{J0..J1\}$
assumes $line-in: \bigwedge x. x \in \{0..1\} \Longrightarrow b + x *R (a - b) \in D$
shows $f a \in \{f b + J0..f b + J1\}$

<proof>

```

end
theory
  Vector-Derivative-On
imports
  HOL-Analysis.Analysis
begin

```

1.24 Vector derivative on a set

definition

```

has-vderiv-on :: (real ⇒ 'a::real-normed-vector) ⇒ (real ⇒ 'a) ⇒ real set ⇒ bool
(infix ‹(has'-vderiv'-on)› 50)

```

where

```

(f has-vderiv-on f') S ‹⟷› (∀ x ∈ S. (f has-vector-derivative f' x) (at x within S))

```

```

lemma has-vderiv-on-empty[intro, simp]: (f has-vderiv-on f') {}
  ‹proof›

```

lemma has-vderiv-on-subset:

```

assumes (f has-vderiv-on f') S
assumes T ⊆ S
shows (f has-vderiv-on f') T
  ‹proof›

```

lemma has-vderiv-on-compose:

```

assumes (f has-vderiv-on f') (g ' T)
assumes (g has-vderiv-on g') T
shows (f o g has-vderiv-on (λx. g' x *R f' (g x))) T
  ‹proof›

```

lemma has-vderiv-on-open:

```

assumes open T
shows (f has-vderiv-on f') T ‹⟷› (∀ t ∈ T. (f has-vector-derivative f' t) (at t))
  ‹proof›

```

lemma has-vderiv-on-eq-rhs:— TODO: integrate intro derivative-eq-intros

```

(f has-vderiv-on g') T ‹⟹› (∧x. x ∈ T ‹⟹› g' x = f' x) ‹⟹› (f has-vderiv-on f')
T
  ‹proof›

```

lemma [THEN has-vderiv-on-eq-rhs, derivative-intros]:

```

shows has-vderiv-on-id: ((λx. x) has-vderiv-on (λx. 1)) T
  and has-vderiv-on-const: ((λx. c) has-vderiv-on (λx. 0)) T
  ‹proof›

```

lemma [THEN has-vderiv-on-eq-rhs, derivative-intros]:

```

fixes f::real ⇒ 'a::real-normed-vector
assumes (f has-vderiv-on f') T

```

shows *has-vderiv-on-uminus*: $((\lambda x. - f x) \text{ has-vderiv-on } (\lambda x. - f' x)) T$
 ⟨*proof*⟩

lemma [*THEN has-vderiv-on-eq-rhs, derivative-intros*]:
fixes $f g :: \text{real} \Rightarrow 'a :: \text{real-normed-vector}$
assumes $(f \text{ has-vderiv-on } f') T$
assumes $(g \text{ has-vderiv-on } g') T$
shows *has-vderiv-on-add*: $((\lambda x. f x + g x) \text{ has-vderiv-on } (\lambda x. f' x + g' x)) T$
and *has-vderiv-on-diff*: $((\lambda x. f x - g x) \text{ has-vderiv-on } (\lambda x. f' x - g' x)) T$
 ⟨*proof*⟩

lemma [*THEN has-vderiv-on-eq-rhs, derivative-intros*]:
fixes $f :: \text{real} \Rightarrow \text{real}$ **and** $g :: \text{real} \Rightarrow 'a :: \text{real-normed-vector}$
assumes $(f \text{ has-vderiv-on } f') T$
assumes $(g \text{ has-vderiv-on } g') T$
shows *has-vderiv-on-scaleR*: $((\lambda x. f x *_{\mathbb{R}} g x) \text{ has-vderiv-on } (\lambda x. f x *_{\mathbb{R}} g' x + f' x *_{\mathbb{R}} g x)) T$
 ⟨*proof*⟩

lemma [*THEN has-vderiv-on-eq-rhs, derivative-intros*]:
fixes $f g :: \text{real} \Rightarrow 'a :: \text{real-normed-algebra}$
assumes $(f \text{ has-vderiv-on } f') T$
assumes $(g \text{ has-vderiv-on } g') T$
shows *has-vderiv-on-mult*: $((\lambda x. f x * g x) \text{ has-vderiv-on } (\lambda x. f x * g' x + f' x * g x)) T$
 ⟨*proof*⟩

lemma *has-vderiv-on-ln* [*THEN has-vderiv-on-eq-rhs, derivative-intros*]:
fixes $g :: \text{real} \Rightarrow \text{real}$
assumes $\bigwedge x. x \in s \implies 0 < g x$
assumes $(g \text{ has-vderiv-on } g') s$
shows $((\lambda x. \ln (g x)) \text{ has-vderiv-on } (\lambda x. g' x / g x)) s$
 ⟨*proof*⟩

lemma *fundamental-theorem-of-calculus'*:
fixes $f :: \text{real} \Rightarrow 'a :: \text{banach}$
shows $a \leq b \implies (f \text{ has-vderiv-on } f') \{a .. b\} \implies (f' \text{ has-integral } (f b - f a)) \{a .. b\}$
 ⟨*proof*⟩

lemma *has-vderiv-on-If*:
assumes $U = S \cup T$
assumes $(f \text{ has-vderiv-on } f') (S \cup (\text{closure } T \cap \text{closure } S))$
assumes $(g \text{ has-vderiv-on } g') (T \cup (\text{closure } T \cap \text{closure } S))$
assumes $\bigwedge x. x \in \text{closure } T \implies x \in \text{closure } S \implies f x = g x$
assumes $\bigwedge x. x \in \text{closure } T \implies x \in \text{closure } S \implies f' x = g' x$
shows $((\lambda t. \text{if } t \in S \text{ then } f t \text{ else } g t) \text{ has-vderiv-on } (\lambda t. \text{if } t \in S \text{ then } f' t \text{ else } g' t)) U$

<proof>

lemma *mut-very-simple-closed-segmentE*:

fixes $f::\text{real}\Rightarrow\text{real}$

assumes $(f \text{ has-vderiv-on } f')$ $(\text{closed-segment } a \ b)$

obtains y **where** $y \in \text{closed-segment } a \ b$ $f \ b - f \ a = (b - a) * f' \ y$

<proof>

lemma *mut-simple-closed-segmentE*:

fixes $f::\text{real}\Rightarrow\text{real}$

assumes $(f \text{ has-vderiv-on } f')$ $(\text{closed-segment } a \ b)$

assumes $a \neq b$

obtains y **where** $y \in \text{open-segment } a \ b$ $f \ b - f \ a = (b - a) * f' \ y$

<proof>

lemma *differentiable-bound-general-open-segment*:

fixes $a :: \text{real}$

and $b :: \text{real}$

and $f :: \text{real} \Rightarrow 'a::\text{real-normed-vector}$

and $f' :: \text{real} \Rightarrow 'a$

assumes $\text{continuous-on } (\text{closed-segment } a \ b)$ f

assumes $\text{continuous-on } (\text{closed-segment } a \ b)$ g

and $(f \text{ has-vderiv-on } f')$ $(\text{open-segment } a \ b)$

and $(g \text{ has-vderiv-on } g')$ $(\text{open-segment } a \ b)$

and $\bigwedge x. x \in \text{open-segment } a \ b \implies \text{norm } (f' \ x) \leq g' \ x$

shows $\text{norm } (f \ b - f \ a) \leq \text{abs } (g \ b - g \ a)$

<proof>

lemma *has-vderiv-on-union*:

assumes $(f \text{ has-vderiv-on } g)$ $(s \cup \text{closure } s \cap \text{closure } t)$

assumes $(f \text{ has-vderiv-on } g)$ $(t \cup \text{closure } s \cap \text{closure } t)$

shows $(f \text{ has-vderiv-on } g)$ $(s \cup t)$

<proof>

lemma *has-vderiv-on-union-closed*:

assumes $(f \text{ has-vderiv-on } g)$ s

assumes $(f \text{ has-vderiv-on } g)$ t

assumes $\text{closed } s$ $\text{closed } t$

shows $(f \text{ has-vderiv-on } g)$ $(s \cup t)$

<proof>

lemma *vderiv-on-continuous-on*: $(f \text{ has-vderiv-on } f') \ S \implies \text{continuous-on } S \ f$

<proof>

lemma *has-vderiv-on-cong[cong]*:

assumes $\bigwedge x. x \in S \implies f \ x = g \ x$

assumes $\bigwedge x. x \in S \implies f' \ x = g' \ x$

assumes $S = T$

shows $(f \text{ has-vderiv-on } f') \ S = (g \text{ has-vderiv-on } g') \ T$

<proof>

lemma *has-vderiv-eq:*

assumes $(f \text{ has-vderiv-on } f')$ S
assumes $\bigwedge x. x \in S \implies f x = g x$
assumes $\bigwedge x. x \in S \implies f' x = g' x$
assumes $S = T$
shows $(g \text{ has-vderiv-on } g')$ T
<proof>

lemma *has-vderiv-on-compose':*

assumes $(f \text{ has-vderiv-on } f')$ $(g \text{ ' } T)$
assumes $(g \text{ has-vderiv-on } g')$ T
shows $((\lambda x. f (g x)) \text{ has-vderiv-on } (\lambda x. g' x *_R f' (g x))) T$
<proof>

lemma *has-vderiv-on-compose2:*

assumes $(f \text{ has-vderiv-on } f')$ S
assumes $(g \text{ has-vderiv-on } g')$ T
assumes $\bigwedge t. t \in T \implies g t \in S$
shows $((\lambda x. f (g x)) \text{ has-vderiv-on } (\lambda x. g' x *_R f' (g x))) T$
<proof>

lemma *has-vderiv-on-singleton:* $(y \text{ has-vderiv-on } y') \{t0\}$

<proof>

lemma

has-vderiv-on-zero-constant:
assumes *convex* s
assumes $(f \text{ has-vderiv-on } (\lambda h. 0)) s$
obtains c **where** $\bigwedge x. x \in s \implies f x = c$
<proof>

lemma *bounded-vderiv-on-imp-lipschitz:*

assumes $(f \text{ has-vderiv-on } f')$ X
assumes *convex:* *convex* X
assumes $\bigwedge x. x \in X \implies \text{norm } (f' x) \leq C \ 0 \leq C$
shows $C\text{-lipschitz-on } X f$
<proof>

end

theory *Interval-Integral-HK*

imports *Vector-Derivative-On*

begin

1.25 interval integral

definition *has-ivl-integral* ::

$(\text{real} \Rightarrow 'b::\text{real-normed-vector}) \Rightarrow 'b \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow \text{bool}$ — TODO: generalize?

(**infixr** $\langle \text{has}'\text{-ivl}'\text{-integral} \rangle$ 46)
where $(f \text{ has-ivl-integral } y) a b \longleftrightarrow (if\ a \leq b\ \text{then}\ (f\ \text{has-integral}\ y)\ \{a\ ..\ b\}\ \text{else}\ (f\ \text{has-integral}\ -\ y)\ \{b\ ..\ a\})$

definition $ivl\text{-integral}::\text{real} \Rightarrow \text{real} \Rightarrow (\text{real} \Rightarrow 'a) \Rightarrow 'a::\text{real-normed-vector}$
where $ivl\text{-integral}\ a\ b\ f = \text{integral}\ \{a\ ..\ b\}\ f - \text{integral}\ \{b\ ..\ a\}\ f$

lemma $integral\text{-emptyI}[simp]$:
fixes $a\ b::\text{real}$
shows $a \geq b \Longrightarrow \text{integral}\ \{a..b\}\ f = 0$ $a > b \Longrightarrow \text{integral}\ \{a..b\}\ f = 0$
 $\langle \text{proof} \rangle$

lemma $ivl\text{-integral-unique}$: $(f \text{ has-ivl-integral } y) a b \Longrightarrow ivl\text{-integral}\ a\ b\ f = y$
 $\langle \text{proof} \rangle$

lemma $fundamental\text{-theorem-of-calculus-ivl-integral}$:
fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$
shows $(f \text{ has-vderiv-on } f')\ (\text{closed-segment}\ a\ b) \Longrightarrow (f' \text{ has-ivl-integral } f\ b - f\ a)$
 $a\ b$
 $\langle \text{proof} \rangle$

lemma
fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$
assumes $f \text{ integrable-on } (\text{closed-segment}\ a\ b)$
shows $indefinite\text{-ivl-integral-continuous}$:
 $continuous\text{-on } (\text{closed-segment}\ a\ b)\ (\lambda x. ivl\text{-integral}\ a\ x\ f)$
 $continuous\text{-on } (\text{closed-segment}\ b\ a)\ (\lambda x. ivl\text{-integral}\ a\ x\ f)$
 $\langle \text{proof} \rangle$

lemma
fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$
assumes $f \text{ integrable-on } (\text{closed-segment}\ a\ b)$
assumes $c \in \text{closed-segment}\ a\ b$
shows $indefinite\text{-ivl-integral-continuous-subset}$:
 $continuous\text{-on } (\text{closed-segment}\ a\ b)\ (\lambda x. ivl\text{-integral}\ c\ x\ f)$
 $\langle \text{proof} \rangle$

lemma $real\text{-Icc-closed-segment}$: **fixes** $a\ b::\text{real}$ **shows** $a \leq b \Longrightarrow \{a\ ..\ b\} = \text{closed-segment}\ a\ b$
 $\langle \text{proof} \rangle$

lemma $ivl\text{-integral-zero}[simp]$: $ivl\text{-integral}\ a\ a\ f = 0$
 $\langle \text{proof} \rangle$

lemma $ivl\text{-integral-cong}$:
assumes $\bigwedge x. x \in \text{closed-segment}\ a\ b \Longrightarrow g\ x = f\ x$
assumes $a = c\ b = d$
shows $ivl\text{-integral}\ a\ b\ f = ivl\text{-integral}\ c\ d\ g$
 $\langle \text{proof} \rangle$

lemma *ivl-integral-diff*:

f integrable-on (closed-segment *s t*) \implies *g* integrable-on (closed-segment *s t*) \implies
ivl-integral s t ($\lambda x. f x - g x$) = *ivl-integral s t* *f* - *ivl-integral s t* *g*
(*proof*)

lemma *ivl-integral-norm-bound-ivl-integral*:

fixes *f* :: real \Rightarrow 'a::banach
assumes *f* integrable-on (closed-segment *a b*)
and *g* integrable-on (closed-segment *a b*)
and $\bigwedge x. x \in \text{closed-segment } a \ b \implies \text{norm } (f x) \leq g x$
shows *norm* (*ivl-integral a b f*) $\leq \text{abs}$ (*ivl-integral a b g*)
(*proof*)

lemma *ivl-integral-norm-bound-integral*:

fixes *f* :: real \Rightarrow 'a::banach
assumes *f* integrable-on (closed-segment *a b*)
and *g* integrable-on (closed-segment *a b*)
and $\bigwedge x. x \in \text{closed-segment } a \ b \implies \text{norm } (f x) \leq g x$
shows *norm* (*ivl-integral a b f*) $\leq \text{integral}$ (closed-segment *a b*) *g*
(*proof*)

lemma *norm-ivl-integral-le*:

fixes *f* :: real \Rightarrow real
assumes *f* integrable-on (closed-segment *a b*)
and *g* integrable-on (closed-segment *a b*)
and $\bigwedge x. x \in \text{closed-segment } a \ b \implies f x \leq g x$
and $\bigwedge x. x \in \text{closed-segment } a \ b \implies 0 \leq f x$
shows *abs* (*ivl-integral a b f*) $\leq \text{abs}$ (*ivl-integral a b g*)
(*proof*)

lemma *ivl-integral-const* [*simp*]:

shows *ivl-integral a b* ($\lambda x. c$) = (*b - a*) *_R *c*
(*proof*)

lemma *ivl-integral-has-vector-derivative*:

fixes *f* :: real \Rightarrow 'a::banach
assumes *continuous-on* (closed-segment *a b*) *f*
and *x* \in closed-segment *a b*
shows (($\lambda u. \text{ivl-integral a u f}$) *has-vector-derivative f x*) (at *x* within closed-segment *a b*)
(*proof*)

lemma *ivl-integral-has-vderiv-on*:

fixes *f* :: real \Rightarrow 'a::banach
assumes *continuous-on* (closed-segment *a b*) *f*
shows (($\lambda u. \text{ivl-integral a u f}$) *has-vderiv-on f*) (closed-segment *a b*)
(*proof*)

lemma *ivl-integral-has-vderiv-on-subset-segment:*

fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$

assumes *continuous-on (closed-segment a b) f*

and $c \in \text{closed-segment } a \ b$

shows $((\lambda u. \text{ivl-integral } c \ u \ f) \text{ has-vderiv-on } f) (\text{closed-segment } a \ b)$

<proof>

lemma *ivl-integral-has-vector-derivative-subset:*

fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$

assumes *continuous-on (closed-segment a b) f*

and $x \in \text{closed-segment } a \ b$

and $c \in \text{closed-segment } a \ b$

shows $((\lambda u. \text{ivl-integral } c \ u \ f) \text{ has-vector-derivative } f \ x) (\text{at } x \ \text{within } \text{closed-segment } a \ b)$

<proof>

lemma

compact-interval-eq-Inf-Sup:

fixes $A::\text{real set}$

assumes *is-interval A compact A A $\neq \{\}$*

shows $A = \{\text{Inf } A \ .. \ \text{Sup } A\}$

<proof>

lemma *ivl-integral-has-vderiv-on-compact-interval:*

fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$

assumes *continuous-on A f*

and $c \in A$ *is-interval A compact A*

shows $((\lambda u. \text{ivl-integral } c \ u \ f) \text{ has-vderiv-on } f) A$

<proof>

lemma *ivl-integral-has-vector-derivative-compact-interval:*

fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$

assumes *continuous-on A f*

and *is-interval A compact A* $x \in A$ $c \in A$

shows $((\lambda u. \text{ivl-integral } c \ u \ f) \text{ has-vector-derivative } f \ x) (\text{at } x \ \text{within } A)$

<proof>

lemma *ivl-integral-combine:*

fixes $f::\text{real} \Rightarrow 'a::\text{banach}$

assumes *f integrable-on (closed-segment a b)*

assumes *f integrable-on (closed-segment b c)*

assumes *f integrable-on (closed-segment a c)*

shows $\text{ivl-integral } a \ b \ f + \text{ivl-integral } b \ c \ f = \text{ivl-integral } a \ c \ f$

<proof>

lemma *integral-equation-swap-initial-value:*

fixes $x::\text{real} \Rightarrow 'a::\text{banach}$

assumes $\bigwedge t. t \in \text{closed-segment } t0 \ t1 \implies x \ t = x \ t0 + \text{ivl-integral } t0 \ t \ (\lambda t. f \ t \ (x \ t))$

assumes $t: t \in \text{closed-segment } t0 \ t1$
assumes $\text{int}: (\lambda t. f \ t \ (x \ t)) \text{ integrable-on closed-segment } t0 \ t1$
shows $x \ t = x \ t1 + \text{ivl-integral } t1 \ t \ (\lambda t. f \ t \ (x \ t))$
 <proof>

lemma *has-integral-nonpos*:
fixes $f :: 'n::\text{euclidean-space} \Rightarrow \text{real}$
assumes $(f \text{ has-integral } i) \ s$
and $\forall x \in s. f \ x \leq 0$
shows $i \leq 0$
 <proof>

lemma *has-ivl-integral-nonneg*:
fixes $f :: \text{real} \Rightarrow \text{real}$
assumes $(f \text{ has-ivl-integral } i) \ a \ b$
and $\bigwedge x. a \leq x \Longrightarrow x \leq b \Longrightarrow 0 \leq f \ x$
and $\bigwedge x. b \leq x \Longrightarrow x \leq a \Longrightarrow f \ x \leq 0$
shows $0 \leq i$
 <proof>

lemma *has-ivl-integral-ivl-integral*:
 $f \text{ integrable-on (closed-segment } a \ b) \longleftrightarrow (f \text{ has-ivl-integral (ivl-integral } a \ b \ f)) \ a \ b$
 <proof>

lemma *ivl-integral-nonneg*:
fixes $f :: \text{real} \Rightarrow \text{real}$
assumes $f \text{ integrable-on (closed-segment } a \ b)$
and $\bigwedge x. a \leq x \Longrightarrow x \leq b \Longrightarrow 0 \leq f \ x$
and $\bigwedge x. b \leq x \Longrightarrow x \leq a \Longrightarrow f \ x \leq 0$
shows $0 \leq \text{ivl-integral } a \ b \ f$
 <proof>

lemma *ivl-integral-bound*:
fixes $f::\text{real} \Rightarrow 'a::\text{banach}$
assumes $\text{continuous-on (closed-segment } a \ b) \ f$
assumes $\bigwedge t. t \in (\text{closed-segment } a \ b) \Longrightarrow \text{norm } (f \ t) \leq B$
shows $\text{norm } (\text{ivl-integral } a \ b \ f) \leq B * \text{abs } (b - a)$
 <proof>

lemma *ivl-integral-minus-sets*:
fixes $f::\text{real} \Rightarrow 'a::\text{banach}$
shows $f \text{ integrable-on (closed-segment } c \ a) \Longrightarrow f \text{ integrable-on (closed-segment } c \ b) \Longrightarrow f \text{ integrable-on (closed-segment } a \ b) \Longrightarrow$
 $\text{ivl-integral } c \ a \ f - \text{ivl-integral } c \ b \ f = \text{ivl-integral } b \ a \ f$
 <proof>

lemma *ivl-integral-minus-sets'*:
fixes $f::\text{real} \Rightarrow 'a::\text{banach}$

shows f integrable-on (closed-segment a c) \implies f integrable-on (closed-segment b c) \implies f integrable-on (closed-segment a b) \implies
 ivl -integral a c f - ivl -integral b c f = ivl -integral a b f
 <proof>

end
theory Gronwall
imports Vector-Derivative-On
begin

1.26 Gronwall

lemma derivative-quotient-bound:

assumes g -deriv-on: (g has-vderiv-on g') $\{a .. b\}$
assumes frac-le: $\bigwedge t. t \in \{a .. b\} \implies g' t / g t \leq K$
assumes g' -cont: continuous-on $\{a .. b\}$ g'
assumes g -pos: $\bigwedge t. t \in \{a .. b\} \implies g t > 0$
assumes t -in: $t \in \{a .. b\}$
shows $g t \leq g a * \exp (K * (t - a))$
 <proof>

lemma derivative-quotient-bound-left:

assumes g -deriv-on: (g has-vderiv-on g') $\{a .. b\}$
assumes frac-ge: $\bigwedge t. t \in \{a .. b\} \implies K \leq g' t / g t$
assumes g' -cont: continuous-on $\{a .. b\}$ g'
assumes g -pos: $\bigwedge t. t \in \{a .. b\} \implies g t > 0$
assumes t -in: $t \in \{a..b\}$
shows $g t \leq g b * \exp (K * (t - b))$
 <proof>

lemma gronwall-general:

fixes g K C a b **and** $t::real$
defines $G \equiv \lambda t. C + K * \text{integral } \{a..t\} (\lambda s. g s)$
assumes g -le- G : $\bigwedge t. t \in \{a..b\} \implies g t \leq G t$
assumes g -cont: continuous-on $\{a..b\}$ g
assumes g -nonneg: $\bigwedge t. t \in \{a..b\} \implies 0 \leq g t$
assumes pos: $0 < C$ $K > 0$
assumes $t \in \{a..b\}$
shows $g t \leq C * \exp (K * (t - a))$
 <proof>

lemma gronwall-general-left:

fixes g K C a b **and** $t::real$
defines $G \equiv \lambda t. C + K * \text{integral } \{t..b\} (\lambda s. g s)$
assumes g -le- G : $\bigwedge t. t \in \{a..b\} \implies g t \leq G t$
assumes g -cont: continuous-on $\{a..b\}$ g
assumes g -nonneg: $\bigwedge t. t \in \{a..b\} \implies 0 \leq g t$
assumes pos: $0 < C$ $K > 0$
assumes $t \in \{a..b\}$

shows $g\ t \leq C * \exp(-K * (t - b))$
 ⟨proof⟩

lemma *gronwall-general-segment*:

fixes $a\ b::\text{real}$
assumes $\bigwedge t. t \in \text{closed-segment } a\ b \implies g\ t \leq C + K * \text{integral } (\text{closed-segment } a\ t)\ g$
and *continuous-on* $(\text{closed-segment } a\ b)\ g$
and $\bigwedge t. t \in \text{closed-segment } a\ b \implies 0 \leq g\ t$
and $0 < C$
and $0 < K$
and $t \in \text{closed-segment } a\ b$
shows $g\ t \leq C * \exp(K * \text{abs}(t - a))$
 ⟨proof⟩

lemma *gronwall-more-general-segment*:

fixes $a\ b\ c::\text{real}$
assumes $\bigwedge t. t \in \text{closed-segment } a\ b \implies g\ t \leq C + K * \text{integral } (\text{closed-segment } c\ t)\ g$
and *cont*: *continuous-on* $(\text{closed-segment } a\ b)\ g$
and $\bigwedge t. t \in \text{closed-segment } a\ b \implies 0 \leq g\ t$
and $0 < C$
and $0 < K$
and $t: t \in \text{closed-segment } a\ b$
and $c: c \in \text{closed-segment } a\ b$
shows $g\ t \leq C * \exp(K * \text{abs}(t - c))$
 ⟨proof⟩

lemma *gronwall*:

fixes $g\ K\ C$ **and** $t::\text{real}$
defines $G \equiv \lambda t. C + K * \text{integral } \{0..t\} (\lambda s. g\ s)$
assumes *g-le-G*: $\bigwedge t. 0 \leq t \implies t \leq a \implies g\ t \leq G\ t$
assumes *g-cont*: *continuous-on* $\{0..a\}\ g$
assumes *g-nonneg*: $\bigwedge t. 0 \leq t \implies t \leq a \implies 0 \leq g\ t$
assumes *pos*: $0 < C\ 0 < K$
assumes $0 \leq t\ t \leq a$
shows $g\ t \leq C * \exp(K * t)$
 ⟨proof⟩

lemma *gronwall-left*:

fixes $g\ K\ C$ **and** $t::\text{real}$
defines $G \equiv \lambda t. C + K * \text{integral } \{t..0\} (\lambda s. g\ s)$
assumes *g-le-G*: $\bigwedge t. a \leq t \implies t \leq 0 \implies g\ t \leq G\ t$
assumes *g-cont*: *continuous-on* $\{a..0\}\ g$
assumes *g-nonneg*: $\bigwedge t. a \leq t \implies t \leq 0 \implies 0 \leq g\ t$
assumes *pos*: $0 < C\ 0 < K$
assumes $a \leq t\ t \leq 0$
shows $g\ t \leq C * \exp(-K * t)$
 ⟨proof⟩

end

2 Initial Value Problems

theory *Initial-Value-Problem*

imports

../ODE-Auxiliarities

../Library/Interval-Integral-HK

../Library/Gronwall

begin

lemma *clamp-le[simp]*: $x \leq a \implies \text{clamp } a \ b \ x = a$ **for** $x::'a::\text{ordered-euclidean-space}$
<proof>

lemma *clamp-ge[simp]*: $a \leq b \implies b \leq x \implies \text{clamp } a \ b \ x = b$ **for** $x::'a::\text{ordered-euclidean-space}$
<proof>

abbreviation *cfuncset* :: $'a::\text{topological-space set} \implies 'b::\text{metric-space set} \implies ('a \implies_C$
 $'b)$ *set*

(**infixr** $\langle \rightarrow_C \rangle$ 60)

where $A \rightarrow_C B \equiv \text{PiC } A \ (\lambda \cdot. B)$

lemma *closed-segment-translation-zero*: $z \in \{z + a \ -- \ z + b\} \longleftrightarrow 0 \in \{a \ -- \ b\}$
<proof>

lemma *closed-segment-subset-interval*: $\text{is-interval } T \implies a \in T \implies b \in T \implies$
 $\text{closed-segment } a \ b \subseteq T$
<proof>

definition *half-open-segment*:: $'a::\text{real-vector} \implies 'a \implies 'a \text{ set } (\langle (1\{\text{---}\langle-\}) \rangle)$
where $\text{half-open-segment } a \ b = \{a \ -- \ b\} - \{b\}$

lemma *half-open-segment-real*:

fixes $a \ b::\text{real}$

shows $\{a \ \text{---}\langle b\} = (\text{if } a \leq b \text{ then } \{a \ ..\langle b\} \text{ else } \{b \ \text{<.. } a\})$

<proof>

lemma *closure-half-open-segment*:

fixes $a \ b::\text{real}$

shows $\text{closure } \{a \ \text{---}\langle b\} = (\text{if } a = b \text{ then } \{\} \text{ else } \{a \ -- \ b\})$

<proof>

lemma *half-open-segment-subset[intro, simp]*:

$\{t0 \ \text{---}\langle t1\} \subseteq \{t0 \ -- \ t1\}$

$x \in \{t0 \ \text{---}\langle t1\} \implies x \in \{t0 \ -- \ t1\}$

<proof>

lemma *half-open-segment-closed-segmentI*:

$t \in \{t0 \text{ -- } t1\} \implies t \neq t1 \implies t \in \{t0 \text{ --} < t1\}$
 ⟨proof⟩

lemma *islimpt-half-open-segment*:

fixes $t0\ t1\ s::real$
assumes $t0 \neq t1\ s \in \{t0 \text{ -- } t1\}$
shows $s \text{ islimpt } \{t0 \text{ --} < t1\}$
 ⟨proof⟩

lemma

mem-half-open-segment-eventually-in-closed-segment:

fixes $t::real$
assumes $t \in \{t0 \text{ --} < t1\}$
shows $\forall_F\ t1' \text{ in at } t1' \text{ within } \{t0 \text{ --} < t1'\}. t \in \{t0 \text{ -- } t1'\}$
 ⟨proof⟩

lemma *closed-segment-half-open-segment-subsetI*:

fixes $x::real$ **shows** $x \in \{t0 \text{ --} < t1\} \implies \{t0 \text{ -- } x\} \subseteq \{t0 \text{ --} < t1\}$
 ⟨proof⟩

lemma *dist-component-le*:

fixes $x\ y::'a::euclidean-space$
assumes $i \in \text{Basis}$
shows $\text{dist } (x \cdot i) (y \cdot i) \leq \text{dist } x\ y$
 ⟨proof⟩

lemma *sum-inner-Basis-one*: $i \in \text{Basis} \implies (\sum_{x \in \text{Basis}} x \cdot i) = 1$

⟨proof⟩

lemma *cball-in-cbox*:

fixes $y::'a::euclidean-space$
shows $\text{cball } y\ r \subseteq \text{cbox } (y - r *_{\mathbb{R}} \text{One}) (y + r *_{\mathbb{R}} \text{One})$
 ⟨proof⟩

lemma *centered-cbox-in-cball*:

shows $\text{cbox } (- r *_{\mathbb{R}} \text{One}) (r *_{\mathbb{R}} \text{One}) \subseteq \text{cball } 0 (\text{sqr}t(\text{DIM } ('a)) * r)$
 ⟨proof⟩

2.1 Solutions of IVPs

definition

solves-ode :: $(real \Rightarrow 'a::real-normed-vector) \Rightarrow (real \Rightarrow 'a \Rightarrow 'a) \Rightarrow real\ set \Rightarrow 'a\ set \Rightarrow bool$

(**infix** ⟨(solves'-ode)⟩ 50)

where

$(y \text{ solves-ode } f) T\ X \iff (y \text{ has-vderiv-on } (\lambda t. f\ t\ (y\ t))) T \wedge y \in T \rightarrow X$

lemma *solves-odeI*:

assumes *solves-ode-vderivD*: $(y \text{ has-vderiv-on } (\lambda t. f t (y t))) T$
and *solves-ode-domainD*: $\bigwedge t. t \in T \implies y t \in X$
shows $(y \text{ solves-ode } f) T X$
 $\langle \text{proof} \rangle$

lemma *solves-odeD*:

assumes $(y \text{ solves-ode } f) T X$
shows *solves-ode-vderivD*: $(y \text{ has-vderiv-on } (\lambda t. f t (y t))) T$
and *solves-ode-domainD*: $\bigwedge t. t \in T \implies y t \in X$
 $\langle \text{proof} \rangle$

lemma *solves-ode-continuous-on*: $(y \text{ solves-ode } f) T X \implies \text{continuous-on } T y$
 $\langle \text{proof} \rangle$

lemma *solves-ode-congI*:

assumes $(x \text{ solves-ode } f) T X$
assumes $\bigwedge t. t \in T \implies x t = y t$
assumes $\bigwedge t. t \in T \implies f t (x t) = g t (x t)$
assumes $T = S X = Y$
shows $(y \text{ solves-ode } g) S Y$
 $\langle \text{proof} \rangle$

lemma *solves-ode-cong[cong]*:

assumes $\bigwedge t. t \in T \implies x t = y t$
assumes $\bigwedge t. t \in T \implies f t (x t) = g t (x t)$
assumes $T = S X = Y$
shows $(x \text{ solves-ode } f) T X \longleftrightarrow (y \text{ solves-ode } g) S Y$
 $\langle \text{proof} \rangle$

lemma *solves-ode-on-subset*:

assumes $(x \text{ solves-ode } f) S Y$
assumes $T \subseteq S Y \subseteq X$
shows $(x \text{ solves-ode } f) T X$
 $\langle \text{proof} \rangle$

lemma *preflect-solution*:

assumes $t0 \in T$
assumes *sol*: $((\lambda t. x (\text{preflect } t0 t)) \text{ solves-ode } (\lambda t x. - f (\text{preflect } t0 t) x))$
 $(\text{preflect } t0 \text{ ' } T) X$
shows $(x \text{ solves-ode } f) T X$
 $\langle \text{proof} \rangle$

lemma *solution-preflect*:

assumes $t0 \in T$
assumes *sol*: $(x \text{ solves-ode } f) T X$
shows $((\lambda t. x (\text{preflect } t0 t)) \text{ solves-ode } (\lambda t x. - f (\text{preflect } t0 t) x)) (\text{preflect } t0$
 $\text{' } T) X$
 $\langle \text{proof} \rangle$

lemma *solution- eq -preflect-solution:*

assumes $t0 \in T$

shows $(x \text{ solves-ode } f) T X \longleftrightarrow ((\lambda t. x (\text{preflect } t0 t)) \text{ solves-ode } (\lambda t x. - f (\text{preflect } t0 t) x)) (\text{preflect } t0 ' T) X$
 $\langle \text{proof} \rangle$

lemma *shift-autonomous-solution:*

assumes $\text{sol}: (x \text{ solves-ode } f) T X$

assumes $\text{auto}: \bigwedge s t. s \in T \implies f s (x s) = f t (x s)$

shows $((\lambda t. x (t + t0)) \text{ solves-ode } f) ((\lambda t. t - t0) ' T) X$

$\langle \text{proof} \rangle$

lemma *solves-ode-singleton:* $y t0 \in X \implies (y \text{ solves-ode } f) \{t0\} X$

$\langle \text{proof} \rangle$

2.1.1 Connecting solutions

lemma *connection-solves-ode:*

assumes $x: (x \text{ solves-ode } f) T X$

assumes $y: (y \text{ solves-ode } g) S Y$

assumes $\text{conn-T}: \text{closure } S \cap \text{closure } T \subseteq T$

assumes $\text{conn-S}: \text{closure } S \cap \text{closure } T \subseteq S$

assumes $\text{conn-x}: \bigwedge t. t \in \text{closure } S \implies t \in \text{closure } T \implies x t = y t$

assumes $\text{conn-f}: \bigwedge t. t \in \text{closure } S \implies t \in \text{closure } T \implies f t (y t) = g t (y t)$

shows $((\lambda t. \text{if } t \in T \text{ then } x t \text{ else } y t) \text{ solves-ode } (\lambda t. \text{if } t \in T \text{ then } f t \text{ else } g t)) (T \cup S) (X \cup Y)$

$\langle \text{proof} \rangle$

lemma

solves-ode-subset-range:

assumes $x: (x \text{ solves-ode } f) T X$

assumes $s: x ' T \subseteq Y$

shows $(x \text{ solves-ode } f) T Y$

$\langle \text{proof} \rangle$

2.2 unique solution with initial value

definition

$\text{usolves-ode-from} :: (\text{real} \Rightarrow 'a::\text{real-normed-vector}) \Rightarrow (\text{real} \Rightarrow 'a \Rightarrow 'a) \Rightarrow \text{real} \Rightarrow \text{real set} \Rightarrow 'a \text{ set} \Rightarrow \text{bool}$

$(\langle (-) \text{ usolves'-ode } (-) \text{ from } (-) \rangle [10, 10, 10] 10)$

— TODO: no idea about mixfix and precedences, check this!

where

$(y \text{ usolves-ode } f \text{ from } t0) T X \longleftrightarrow (y \text{ solves-ode } f) T X \wedge t0 \in T \wedge \text{is-interval } T \wedge$

$(\forall z T'. t0 \in T' \wedge \text{is-interval } T' \wedge T' \subseteq T \wedge (z \text{ solves-ode } f) T' X \longrightarrow z t0 = y t0 \longrightarrow (\forall t \in T'. z t = y t))$

uniqueness of solution can depend on domain X :

lemma

$((\lambda-. 0::real) \text{ usolves-ode } (\lambda-. \text{sqrt}) \text{ from } 0) \{0..\} \{0\}$
 $((\lambda t. t^2 / 4) \text{ solves-ode } (\lambda-. \text{sqrt})) \{0..\} \{0..\}$
 $(\lambda t. t^2 / 4) 0 = (\lambda-. 0::real) 0$
 $\langle \text{proof} \rangle$

TODO: show that if solution stays in interior, then domain can be enlarged!
(?)

lemma *usolves-odeD*:

assumes $(y \text{ usolves-ode } f \text{ from } t0) T X$
shows $(y \text{ solves-ode } f) T X$
and $t0 \in T$
and *is-interval* T
and $\bigwedge z T' t. t0 \in T' \implies \text{is-interval } T' \implies T' \subseteq T \implies (z \text{ solves-ode } f) T' X$
 $\implies z t0 = y t0 \implies t \in T' \implies z t = y t$
 $\langle \text{proof} \rangle$

lemma *usolves-ode-rawI*:

assumes $(y \text{ solves-ode } f) T X t0 \in T \text{ is-interval } T$
assumes $\bigwedge z T' t. t0 \in T' \implies \text{is-interval } T' \implies T' \subseteq T \implies (z \text{ solves-ode } f)$
 $T' X \implies z t0 = y t0 \implies t \in T' \implies z t = y t$
shows $(y \text{ usolves-ode } f \text{ from } t0) T X$
 $\langle \text{proof} \rangle$

lemma *usolves-odeI*:

assumes $(y \text{ solves-ode } f) T X t0 \in T \text{ is-interval } T$
assumes *usol*: $\bigwedge z t. \{t0 \text{ -- } t\} \subseteq T \implies (z \text{ solves-ode } f) \{t0 \text{ -- } t\} X \implies z t0$
 $= y t0 \implies z t = y t$
shows $(y \text{ usolves-ode } f \text{ from } t0) T X$
 $\langle \text{proof} \rangle$

lemma *is-interval-singleton*[*intro,simp*]: *is-interval* $\{t0\}$

$\langle \text{proof} \rangle$

lemma *usolves-ode-singleton*: $x t0 \in X \implies (x \text{ usolves-ode } f \text{ from } t0) \{t0\} X$

$\langle \text{proof} \rangle$

lemma *usolves-ode-congI*:

assumes $x: (x \text{ usolves-ode } f \text{ from } t0) T X$
assumes $\bigwedge t. t \in T \implies x t = y t$
assumes $\bigwedge t y. t \in T \implies y \in X \implies f t y = g t y$ — TODO: weaken this assumption?!
assumes $t0 = s0$
assumes $T = S$
assumes $X = Y$
shows $(y \text{ usolves-ode } g \text{ from } s0) S Y$
 $\langle \text{proof} \rangle$

lemma *usolves-ode-cong[cong]*:

assumes $\bigwedge t. t \in T \implies x t = y t$
assumes $\bigwedge t y. t \in T \implies y \in X \implies f t y = g t y$ — TODO: weaken this assumption?!
assumes $t0 = s0$
assumes $T = S$
assumes $X = Y$
shows $(x \text{ usolves-ode } f \text{ from } t0) T X \longleftrightarrow (y \text{ usolves-ode } g \text{ from } s0) S Y$
 $\langle \text{proof} \rangle$

lemma *shift-autonomous-unique-solution*:

assumes *usol*: $(x \text{ usolves-ode } f \text{ from } t0) T X$
assumes *auto*: $\bigwedge s t x. x \in X \implies f s x = f t x$
shows $((\lambda t. x (t + t0 - t1)) \text{ usolves-ode } f \text{ from } t1) ((+) (t1 - t0) ' T) X$
 $\langle \text{proof} \rangle$

lemma *three-intervals-lemma*:

fixes $a b c :: \text{real}$
assumes $a: a \in A - B$
and $b: b \in B - A$
and $c: c \in A \cap B$
and *iA*: *is-interval* A **and** *iB*: *is-interval* B
and *aI*: $a \in I$
and *bI*: $b \in I$
and *iI*: *is-interval* I
shows $c \in I$
 $\langle \text{proof} \rangle$

lemma *connection-usolves-ode*:

assumes *x*: $(x \text{ usolves-ode } f \text{ from } tx) T X$
assumes *y*: $\bigwedge t. t \in \text{closure } S \cap \text{closure } T \implies (y \text{ usolves-ode } g \text{ from } t) S X$
assumes *conn-T*: $\text{closure } S \cap \text{closure } T \subseteq T$
assumes *conn-S*: $\text{closure } S \cap \text{closure } T \subseteq S$
assumes *conn-t*: $t \in \text{closure } S \cap \text{closure } T$
assumes *conn-x*: $\bigwedge t. t \in \text{closure } S \implies t \in \text{closure } T \implies x t = y t$
assumes *conn-f*: $\bigwedge t x. t \in \text{closure } S \implies t \in \text{closure } T \implies x \in X \implies f t x = g t x$
shows $((\lambda t. \text{if } t \in T \text{ then } x t \text{ else } y t) \text{ usolves-ode } (\lambda t. \text{if } t \in T \text{ then } f t \text{ else } g t) \text{ from } tx) (T \cup S) X$
 $\langle \text{proof} \rangle$

lemma *usolves-ode-union-closed*:

assumes *x*: $(x \text{ usolves-ode } f \text{ from } tx) T X$
assumes *y*: $\bigwedge t. t \in \text{closure } S \cap \text{closure } T \implies (x \text{ usolves-ode } f \text{ from } t) S X$
assumes *conn-T*: $\text{closure } S \cap \text{closure } T \subseteq T$
assumes *conn-S*: $\text{closure } S \cap \text{closure } T \subseteq S$
assumes *conn-t*: $t \in \text{closure } S \cap \text{closure } T$
shows $(x \text{ usolves-ode } f \text{ from } tx) (T \cup S) X$
 $\langle \text{proof} \rangle$

lemma *usolves-ode-solves-odeI*:
assumes $(x \text{ usolves-ode } f \text{ from } tx) \ T \ X$
assumes $(y \text{ solves-ode } f) \ T \ X \ y \ tx = x \ tx$
shows $(y \text{ usolves-ode } f \text{ from } tx) \ T \ X$
 $\langle \text{proof} \rangle$

lemma *usolves-ode-subset-range*:
assumes $x: (x \text{ usolves-ode } f \text{ from } t0) \ T \ X$
assumes $r: x \ ' \ T \subseteq Y \ \mathbf{and} \ Y \subseteq X$
shows $(x \text{ usolves-ode } f \text{ from } t0) \ T \ Y$
 $\langle \text{proof} \rangle$

2.3 ivp on interval

context
fixes $t0 \ t1::\text{real} \ \mathbf{and} \ T$
defines $T \equiv \text{closed-segment } t0 \ t1$
begin

lemma *is-solution-ext-cont*:
 $\text{continuous-on } T \ x \implies (\text{ext-cont } x \ (\text{min } t0 \ t1) \ (\text{max } t0 \ t1) \ \text{solves-ode } f) \ T \ X =$
 $(x \text{ solves-ode } f) \ T \ X$
 $\langle \text{proof} \rangle$

lemma *solution-fixed-point*:
fixes $x::\text{real} \Rightarrow 'a::\text{banach}$
assumes $x: (x \text{ solves-ode } f) \ T \ X \ \mathbf{and} \ t: t \in T$
shows $x \ t0 + \text{ivl-integral } t0 \ t \ (\lambda t. f \ t \ (x \ t)) = x \ t$
 $\langle \text{proof} \rangle$

lemma *solution-fixed-point-left*:
fixes $x::\text{real} \Rightarrow 'a::\text{banach}$
assumes $x: (x \text{ solves-ode } f) \ T \ X \ \mathbf{and} \ t: t \in T$
shows $x \ t1 - \text{ivl-integral } t \ t1 \ (\lambda t. f \ t \ (x \ t)) = x \ t$
 $\langle \text{proof} \rangle$

lemma *solution-fixed-pointI*:
fixes $x::\text{real} \Rightarrow 'a::\text{banach}$
assumes $\text{cont-f: continuous-on } (T \times X) \ (\lambda(t, x). f \ t \ x)$
assumes $\text{cont-x: continuous-on } T \ x$
assumes $\text{defined: } \bigwedge t. t \in T \implies x \ t \in X$
assumes $\text{fp: } \bigwedge t. t \in T \implies x \ t = x \ t0 + \text{ivl-integral } t0 \ t \ (\lambda t. f \ t \ (x \ t))$
shows $(x \text{ solves-ode } f) \ T \ X$
 $\langle \text{proof} \rangle$

end

lemma *solves-ode-half-open-segment-continuation*:

```

fixes  $f::real \Rightarrow 'a \Rightarrow 'a::banach$ 
assumes  $ode: (x \text{ solves-ode } f) \{t0 \text{ --< } t1\} X$ 
assumes  $continuous: continuous-on (\{t0 \text{ -- } t1\} \times X) (\lambda(t, x). f t x)$ 
assumes  $compact X$ 
assumes  $t0 \neq t1$ 
obtains  $l$  where
   $(x \longrightarrow l)$   $(at\ t1\ \text{within } \{t0 \text{ --< } t1\})$ 
   $((\lambda t. \text{ if } t = t1 \text{ then } l \text{ else } x\ t) \text{ solves-ode } f) \{t0 \text{ -- } t1\} X$ 
 $\langle proof \rangle$ 

```

2.4 Picard-Lindelof on set of functions into closed set

```

locale  $continuous-rhs = \text{fixes } T\ X\ f$ 
assumes  $continuous: continuous-on (T \times X) (\lambda(t, x). f t x)$ 
begin

```

```

lemma  $continuous-rhs-comp[continuous-intros]:$ 
assumes  $[continuous-intros]: continuous-on\ S\ g$ 
assumes  $[continuous-intros]: continuous-on\ S\ h$ 
assumes  $g \text{ ' } S \subseteq T$ 
assumes  $h \text{ ' } S \subseteq X$ 
shows  $continuous-on\ S (\lambda x. f (g\ x) (h\ x))$ 
 $\langle proof \rangle$ 

```

end

```

locale  $global-lipschitz =$ 
fixes  $T\ X\ f$  and  $L::real$ 
assumes  $lipschitz: \bigwedge t. t \in T \implies L\text{-lipschitz-on } X (\lambda x. f t x)$ 

```

```

locale  $closed-domain =$ 
fixes  $X$  assumes  $closed: closed\ X$ 

```

```

locale  $interval = \text{fixes } T::real\ \text{set}$ 
assumes  $interval: is\text{-interval } T$ 
begin

```

```

lemma  $closed-segment-subset-domain: t0 \in T \implies t \in T \implies closed\text{-segment } t0\ t$ 
 $\subseteq T$ 
 $\langle proof \rangle$ 

```

```

lemma  $closed-segment-subset-domainI: t0 \in T \implies t \in T \implies s \in closed\text{-segment}$ 
 $t0\ t \implies s \in T$ 
 $\langle proof \rangle$ 

```

```

lemma  $convex[intro, simp]: convex\ T$ 
and  $connected[intro, simp]: connected\ T$ 
 $\langle proof \rangle$ 

```

end

locale *nonempty-set* = **fixes** T **assumes** *nonempty-set*: $T \neq \{\}$

locale *compact-interval* = *interval* + *nonempty-set* T +
assumes *compact-time*: *compact* T
begin

definition $tmin = \text{Inf } T$

definition $tmax = \text{Sup } T$

lemma

shows $tmin$: $t \in T \implies tmin \leq t$ $tmin \in T$
and $tmax$: $t \in T \implies t \leq tmax$ $tmax \in T$
<proof>

lemma *tmin-le-tmax*[*intro, simp*]: $tmin \leq tmax$
<proof>

lemma *T-def*: $T = \{tmin .. tmax\}$
<proof>

lemma *mem-T-I*[*intro, simp*]: $tmin \leq t \implies t \leq tmax \implies t \in T$
<proof>

end

locale *self-mapping* = *interval* T **for** T +

fixes $t0::\text{real}$ **and** $x0 f X$

assumes *iv-defined*: $t0 \in T$ $x0 \in X$

assumes *self-mapping*:

$\bigwedge x t. t \in T \implies x t0 = x0 \implies x \in \text{closed-segment } t0 t \rightarrow X \implies$

$\text{continuous-on } (\text{closed-segment } t0 t) x \implies x t0 + \text{ivl-integral } t0 t (\lambda t. f t (x$
 $t)) \in X$

begin

sublocale *nonempty-set* T *<proof>*

lemma *closed-segment-iv-subset-domain*: $t \in T \implies \text{closed-segment } t0 t \subseteq T$
<proof>

end

locale *unique-on-closed* =

compact-interval T +

self-mapping T $t0$ $x0 f X$ +

continuous-rhs $T X f$ +

closed-domain X +

global-lipschitz $T X f L$ **for** $t0::\text{real}$ **and** T **and** $x0::'a::\text{banach}$ **and** $f X L$

begin

lemma *T-split*: $T = \{tmin .. t0\} \cup \{t0 .. tmax\}$
<proof>

lemma *L-nonneg*: $0 \leq L$
<proof>

Picard Iteration

definition *P-inner* **where** *P-inner* $x\ t = x0 + ivl\text{-integral}\ t0\ t\ (\lambda t. f\ t\ (x\ t))$

definition *P*::($real \Rightarrow_C 'a$) \Rightarrow ($real \Rightarrow_C 'a$)
where *P* $x = (SOME\ g::real \Rightarrow_C 'a.$
 $(\forall t \in T. g\ t = P\text{-inner}\ x\ t) \wedge$
 $(\forall t \leq tmin. g\ t = P\text{-inner}\ x\ tmin) \wedge$
 $(\forall t \geq tmax. g\ t = P\text{-inner}\ x\ tmax))$

lemma *cont-P-inner-ivl*:
 $x \in T \rightarrow_C X \Longrightarrow \text{continuous-on}\ \{tmin..tmax\}\ (P\text{-inner}\ (\text{apply-bcontfun}\ x))$
<proof>

lemma *P-inner-t0[simp]*: $P\text{-inner}\ g\ t0 = x0$
<proof>

lemma *t0-cs-tmin-tmax*: $t0 \in \{tmin--tmax\}$ **and** *cs-tmin-tmax-subset*: $\{tmin--tmax\} \subseteq T$
<proof>

lemma
P-eqs:
assumes $x \in T \rightarrow_C X$
shows *P-eq-P-inner*: $t \in T \Longrightarrow P\ x\ t = P\text{-inner}\ x\ t$
 and *P-le-tmin*: $t \leq tmin \Longrightarrow P\ x\ t = P\text{-inner}\ x\ tmin$
 and *P-ge-tmax*: $t \geq tmax \Longrightarrow P\ x\ t = P\text{-inner}\ x\ tmax$
<proof>

lemma *P-if-eq*:
 $x \in T \rightarrow_C X \Longrightarrow$
 $P\ x\ t = (\text{if}\ tmin \leq t \wedge t \leq tmax\ \text{then}\ P\text{-inner}\ x\ t\ \text{else if}\ t \geq tmax\ \text{then}\ P\text{-inner}\ x\ tmax\ \text{else}\ P\text{-inner}\ x\ tmin)$
<proof>

lemma *dist-P-le*:
assumes $y: y \in T \rightarrow_C X$ **and** $z: z \in T \rightarrow_C X$
assumes *le*: $\bigwedge t. tmin \leq t \Longrightarrow t \leq tmax \Longrightarrow \text{dist}\ (P\text{-inner}\ y\ t)\ (P\text{-inner}\ z\ t) \leq R$
assumes $0 \leq R$
shows $\text{dist}\ (P\ y\ t)\ (P\ z\ t) \leq R$
<proof>

lemma *P-def'*:

assumes $t \in T$

assumes $fixed_point \in T \rightarrow_C X$

shows $(P\ fixed_point)\ t = x0 + ivl_integral\ t0\ t\ (\lambda x. f\ x\ (fixed_point\ x))$

<proof>

definition $iter_space = PiC\ T\ ((\lambda-. X)(t0:=\{x0\}))$

lemma *iter-spaceI*:

assumes $g \in T \rightarrow_C X\ g\ t0 = x0$

shows $g \in iter_space$

<proof>

lemma *iter-spaceD*:

assumes $g \in iter_space$

shows $g \in T \rightarrow_C X\ apply_bcontfun\ g\ t0 = x0$

<proof>

lemma *const-in-iter-space*: $const_bcontfun\ x0 \in iter_space$

<proof>

lemma *closed-iter-space*: $closed\ iter_space$

<proof>

lemma *iter-space-notempty*: $iter_space \neq \{\}$

<proof>

lemma *clamp-in-eq[simp]*: **fixes** $a\ x\ b::real$ **shows** $a \leq x \implies x \leq b \implies clamp\ a$

$b\ x = x$

<proof>

lemma *P-self-mapping*:

assumes $in_space: g \in iter_space$

shows $P\ g \in iter_space$

<proof>

lemma *continuous-on-T*: $continuous_on\ \{tmin..tmax\}\ g \implies continuous_on\ T\ g$

<proof>

lemma *T-closed-segment-subsetI[intro, simp]*: $t \in \{tmin..tmax\} \implies t \in T$

and $T_subsetI[intro, simp]$: $tmin \leq t \implies t \leq tmax \implies t \in T$

<proof>

lemma *t0-mem-closed-segment[intro, simp]*: $t0 \in \{tmin..tmax\}$

<proof>

lemma *tmin-le-t0[intro, simp]*: $tmin \leq t0$

and *tmax-ge-t0[intro, simp]*: $tmax \geq t0$

<proof>

lemma *apply-bcontfun-solution-fixed-point:*

assumes *ode:* (*apply-bcontfun x solves-ode f*) $T X$

assumes *iv:* $x t0 = x0$

assumes *t:* $t \in T$

shows $P x t = x t$

<proof>

lemma

solution-in-iter-space:

assumes *ode:* (*apply-bcontfun z solves-ode f*) $T X$

assumes *iv:* $z t0 = x0$

shows $z \in \text{iter-space}$ (**is** $?z \in -$)

<proof>

end

locale *unique-on-bounded-closed = unique-on-closed +*

assumes *lipschitz-bound:* $\bigwedge s t. s \in T \implies t \in T \implies \text{abs } (s - t) * L < 1$

begin

lemma *lipschitz-bound-maxmin:* $(tmax - tmin) * L < 1$

<proof>

lemma *lipschitz-P:*

shows $((tmax - tmin) * L)$ -*lipschitz-on iter-space P*

<proof>

lemma *fixed-point-unique:* $\exists! x \in \text{iter-space}. P x = x$

<proof>

definition *fixed-point where*

fixed-point = (THE x. x \in iter-space \wedge P x = x)

lemma *fixed-point':*

fixed-point \in iter-space \wedge P fixed-point = fixed-point

<proof>

lemma *fixed-point:*

fixed-point \in iter-space P fixed-point = fixed-point

<proof>

lemma *fixed-point-equality':* $x \in \text{iter-space} \wedge P x = x \implies \text{fixed-point} = x$

<proof>

lemma *fixed-point-equality:* $x \in \text{iter-space} \implies P x = x \implies \text{fixed-point} = x$

<proof>

lemma *fixed-point-iv*: *fixed-point* $t0 = x0$
and *fixed-point-domain*: $x \in T \implies \text{fixed-point } x \in X$
 $\langle \text{proof} \rangle$

lemma *fixed-point-has-vderiv-on*: (*fixed-point has-vderiv-on* ($\lambda t. f t (\text{fixed-point } t)$))
 T
 $\langle \text{proof} \rangle$

lemma *fixed-point-solution*:
shows (*fixed-point solves-ode* f) $T X$
 $\langle \text{proof} \rangle$

2.4.1 Unique solution

lemma *solves-ode-equals-fixed-point*:
assumes *ode*: (*x solves-ode* f) $T X$
assumes *iv*: $x t0 = x0$
assumes *t*: $t \in T$
shows $x t = \text{fixed-point } t$
 $\langle \text{proof} \rangle$

lemma *solves-ode-on-closed-segment-equals-fixed-point*:
assumes *ode*: (*x solves-ode* f) $\{t0 \text{ -- } t1\} X$
assumes *iv*: $x t0 = x0$
assumes *subset*: $\{t0 \text{ -- } t1\} \subseteq T$
assumes *t-mem*: $t \in \{t0 \text{ -- } t1\}$
shows $x t = \text{fixed-point } t$
 $\langle \text{proof} \rangle$

lemma *unique-solution*:
assumes *ivp1*: (*x solves-ode* f) $T X$ $x t0 = x0$
assumes *ivp2*: (*y solves-ode* f) $T X$ $y t0 = x0$
assumes *t*: $t \in T$
shows $x t = y t$
 $\langle \text{proof} \rangle$

lemma *fixed-point-usolves-ode*: (*fixed-point usolves-ode* f from $t0$) $T X$
 $\langle \text{proof} \rangle$

end

lemma *closed-segment-Un*:
fixes $a b c :: \text{real}$
assumes $b \in \text{closed-segment } a c$
shows $\text{closed-segment } a b \cup \text{closed-segment } b c = \text{closed-segment } a c$
 $\langle \text{proof} \rangle$

lemma *closed-segment-closed-segment-subset*:

fixes $s::\text{real}$ **and** $i::\text{nat}$
assumes $s \in \text{closed-segment } a \ b$
assumes $a \in \text{closed-segment } c \ d$ $b \in \text{closed-segment } c \ d$
shows $s \in \text{closed-segment } c \ d$
 $\langle \text{proof} \rangle$

context *unique-on-closed* **begin**

context— solution until $t1$
fixes $t1::\text{real}$
assumes $\text{mem-}t1: t1 \in T$
begin

lemma *subdivide-count-ex*: $\exists n. L * \text{abs } (t1 - t0) / (\text{Suc } n) < 1$
 $\langle \text{proof} \rangle$

definition *subdivide-count* = $(\text{SOME } n. L * \text{abs } (t1 - t0) / \text{Suc } n < 1)$

lemma *subdivide-count*: $L * \text{abs } (t1 - t0) / \text{Suc } \text{subdivide-count} < 1$
 $\langle \text{proof} \rangle$

lemma *subdivide-lipschitz*:
assumes $|s - t| \leq \text{abs } (t1 - t0) / \text{Suc } \text{subdivide-count}$
shows $|s - t| * L < 1$
 $\langle \text{proof} \rangle$

lemma *subdivide-lipschitz-lemma*:
assumes $st: s \in \{a \ \text{--} \ b\}$ $t \in \{a \ \text{--} \ b\}$
assumes $\text{abs } (b - a) \leq \text{abs } (t1 - t0) / \text{Suc } \text{subdivide-count}$
shows $|s - t| * L < 1$
 $\langle \text{proof} \rangle$

definition *step* = $(t1 - t0) / \text{Suc } \text{subdivide-count}$

lemma *last-step*: $t0 + \text{real } (\text{Suc } \text{subdivide-count}) * \text{step} = t1$
 $\langle \text{proof} \rangle$

lemma *step-in-segment*:
assumes $0 \leq i$ $i \leq \text{real } (\text{Suc } \text{subdivide-count})$
shows $t0 + i * \text{step} \in \text{closed-segment } t0 \ t1$
 $\langle \text{proof} \rangle$

lemma *subset-T1*:
fixes $s::\text{real}$ **and** $i::\text{nat}$
assumes $s \in \text{closed-segment } t0 \ (t0 + i * \text{step})$
assumes $i \leq \text{Suc } \text{subdivide-count}$
shows $s \in \{t0 \ \text{--} \ t1\}$
 $\langle \text{proof} \rangle$

lemma *subset-T*: $\{t0 \text{ -- } t1\} \subseteq T$ **and** *subset-TI*: $s \in \{t0 \text{ -- } t1\} \implies s \in T$
 ⟨*proof*⟩

primrec *psolution*:: $\text{nat} \Rightarrow \text{real} \Rightarrow 'a$ **where**
psolution 0 $t = x0$
 | *psolution* (Suc i) $t = \text{unique-on-bounded-closed.fixed-point}$
 $(t0 + \text{real } i * \text{step}) \{t0 + \text{real } i * \text{step} \text{ -- } t0 + \text{real } (\text{Suc } i) * \text{step}\}$
 $(\text{psolution } i (t0 + \text{real } i * \text{step})) f X t$

definition *psolutions* $t = \text{psolution}$ (LEAST $i. t \in \text{closed-segment } (t0 + \text{real } (i - 1) * \text{step}) (t0 + \text{real } i * \text{step})) t$

lemma *psolutions-usolves-until-step*:
assumes *i-le*: $i \leq \text{Suc } \text{subdivide-count}$
shows $(\text{psolutions usolves-ode } f \text{ from } t0) (\text{closed-segment } t0 (t0 + \text{real } i * \text{step}))$
 X
 ⟨*proof*⟩

lemma *psolutions-usolves-ode*: $(\text{psolutions usolves-ode } f \text{ from } t0) \{t0 \text{ -- } t1\} X$
 ⟨*proof*⟩

end

definition *solution* $t = (\text{if } t \leq t0 \text{ then } \text{psolutions } tmin \text{ } t \text{ else } \text{psolutions } tmax \text{ } t)$

lemma *solution-eq-left*: $tmin \leq t \implies t \leq t0 \implies \text{solution } t = \text{psolutions } tmin \text{ } t$
 ⟨*proof*⟩

lemma *solution-eq-right*: $t0 \leq t \implies t \leq tmax \implies \text{solution } t = \text{psolutions } tmax \text{ } t$
 ⟨*proof*⟩

lemma *solution-usolves-ode*: $(\text{solution usolves-ode } f \text{ from } t0) T X$
 ⟨*proof*⟩

lemma *solution-solves-ode*: $(\text{solution solves-ode } f) T X$
 ⟨*proof*⟩

lemma *solution-iv[simp]*: $\text{solution } t0 = x0$
 ⟨*proof*⟩

end

2.5 Picard-Lindelof for $X = UNIV$

locale *unique-on-strip* =
compact-interval $T +$
continuous-rhs $T UNIV f +$
global-lipschitz $T UNIV f L$

```

for  $t0$  and  $T$  and  $f::real \Rightarrow 'a \Rightarrow 'a::banach$  and  $L +$ 
assumes  $iv-time: t0 \in T$ 
begin

sublocale  $unique-on-closed\ t0\ T\ x0\ f\ UNIV\ L$  for  $x0$ 
   $\langle proof \rangle$ 

end

```

2.6 Picard-Lindelof on cylindric domain

```

locale  $solution-in-cylinder =$ 
   $continuous-rhs\ T\ cball\ x0\ b\ f +$ 
   $compact-interval\ T$ 
for  $t0\ T\ x0\ b$  and  $f::real \Rightarrow 'a \Rightarrow 'a::banach +$ 
fixes  $X\ B$ 
defines  $X \equiv cball\ x0\ b$ 
assumes  $initial-time-in: t0 \in T$ 
assumes  $norm-f: \bigwedge x\ t. t \in T \Longrightarrow x \in X \Longrightarrow norm\ (f\ t\ x) \leq B$ 
assumes  $b-pos: b \geq 0$ 
assumes  $e-bounded: \bigwedge t. t \in T \Longrightarrow dist\ t\ t0 \leq b / B$ 
begin

```

```

lemmas  $cylinder = X-def$ 

```

```

lemma  $B-nonneg: B \geq 0$ 
   $\langle proof \rangle$ 

```

```

lemma  $in-bounds-derivativeI:$ 
  assumes  $t \in T$ 
  assumes  $init: x\ t0 = x0$ 
  assumes  $cont: continuous-on\ (closed-segment\ t0\ t)\ x$ 
  assumes  $solves: (x\ has-vderiv-on\ (\lambda s. f\ s\ (y\ s)))\ (open-segment\ t0\ t)$ 
  assumes  $y-bounded: \bigwedge \xi. \xi \in closed-segment\ t0\ t \Longrightarrow x\ \xi \in X \Longrightarrow y\ \xi \in X$ 
  shows  $x\ t \in cball\ x0\ (B * abs\ (t - t0))$ 
   $\langle proof \rangle$ 

```

```

lemma  $in-bounds-derivative-globalI:$ 
  assumes  $t \in T$ 
  assumes  $init: x\ t0 = x0$ 
  assumes  $cont: continuous-on\ (closed-segment\ t0\ t)\ x$ 
  assumes  $solves: (x\ has-vderiv-on\ (\lambda s. f\ s\ (y\ s)))\ (open-segment\ t0\ t)$ 
  assumes  $y-bounded: \bigwedge \xi. \xi \in closed-segment\ t0\ t \Longrightarrow x\ \xi \in X \Longrightarrow y\ \xi \in X$ 
  shows  $x\ t \in X$ 
   $\langle proof \rangle$ 

```

```

lemma  $integral-in-bounds:$ 
  assumes  $t \in T\ x\ t0 = x0\ x \in \{t0 \ --\ t\} \rightarrow X$ 
  assumes  $cont[continuous-intros]: continuous-on\ (\{t0 \ --\ t\})\ x$ 

```

shows $x\ t0 + ivl\text{-integral}\ t0\ t\ (\lambda t. f\ t\ (x\ t)) \in X$ (**is** - + ?ix $t \in X$)
(proof)

lemma *solves-in-cone*:

assumes $t \in T$

assumes *init*: $x\ t0 = x0$

assumes *cont*: *continuous-on* (*closed-segment* $t0\ t$) x

assumes *solves*: (*x has-vderiv-on* ($\lambda s. f\ s\ (x\ s)$)) (*open-segment* $t0\ t$)

shows $x\ t \in cball\ x0\ (B * abs\ (t - t0))$

(proof)

lemma *is-solution-in-cone*:

assumes $t \in T$

assumes *sol*: (*x solves-ode* f) (*closed-segment* $t0\ t$) Y **and** *iv*: $x\ t0 = x0$

shows $x\ t \in cball\ x0\ (B * abs\ (t - t0))$

(proof)

lemma *cone-subset-domain*:

assumes $t \in T$

shows $cball\ x0\ (B * |t - t0|) \subseteq X$

(proof)

lemma *is-solution-in-domain*:

assumes $t \in T$

assumes *sol*: (*x solves-ode* f) (*closed-segment* $t0\ t$) Y **and** *iv*: $x\ t0 = x0$

shows $x\ t \in X$

(proof)

lemma *solves-ode-on-subset-domain*:

assumes *sol*: (*x solves-ode* f) $S\ Y$ **and** *iv*: $x\ t0 = x0$

and *ivl*: $t0 \in S$ *is-interval* $S\ S \subseteq T$

shows (*x solves-ode* f) $S\ X$

(proof)

lemma *usolves-ode-on-subset*:

assumes *x*: (*x usolves-ode* f *from* $t0$) $T\ X$ **and** *iv*: $x\ t0 = x0$

assumes $t0 \in S$ *is-interval* $S\ S \subseteq T\ X \subseteq Y$

shows (*x usolves-ode* f *from* $t0$) $S\ Y$

(proof)

lemma *usolves-ode-on-superset-domain*:

assumes (*x usolves-ode* f *from* $t0$) $T\ X$ **and** *iv*: $x\ t0 = x0$

assumes $X \subseteq Y$

shows (*x usolves-ode* f *from* $t0$) $T\ Y$

(proof)

end

locale *unique-on-cylinder* =

```

    solution-in-cylinder t0 T x0 b f X B +
    global-lipschitz T X f L
    for t0 T x0 b X f B L
begin

sublocale unique-on-closed t0 T x0 f X L
  ⟨proof⟩

end

locale derivative-on-prod =
  fixes T X and f::real ⇒ 'a::banach ⇒ 'a and f':: real × 'a ⇒ (real × 'a) ⇒ 'a
  assumes f': ∧tx. tx ∈ T × X ⇒ ((λ(t, x). f t x) has-derivative (f' tx)) (at tx
  within (T × X))
begin

lemma f'-comp[derivative-intros]:
  (g has-derivative g') (at s within S) ⇒ (h has-derivative h') (at s within S) ⇒
  s ∈ S ⇒ (∧x. x ∈ S ⇒ g x ∈ T) ⇒ (∧x. x ∈ S ⇒ h x ∈ X) ⇒
  ((λx. f (g x) (h x)) has-derivative (λy. f' (g s, h s) (g' y, h' y))) (at s within S)
  ⟨proof⟩

lemma derivative-on-prod-subset:
  assumes X' ⊆ X
  shows derivative-on-prod T X' f f'
  ⟨proof⟩

end

end

theory Picard-Lindelof-Qualitative
imports Initial-Value-Problem
begin

```

2.7 Picard-Lindelof On Open Domains

2.7.1 Local Solution with local Lipschitz

```

lemma cball-eq-closed-segment-real:
  fixes x e::real
  shows cball x e = (if e ≥ 0 then {x - e .. x + e} else {})
  ⟨proof⟩

lemma cube-in-cball:
  fixes x y :: 'a::euclidean-space
  assumes r > 0
  assumes ∧i. i ∈ Basis ⇒ dist (x · i) (y · i) ≤ r / sqrt(DIM('a))
  shows y ∈ cball x r
  ⟨proof⟩

```

lemma *cbox-in-cball'*:
fixes $x::'a::\text{euclidean-space}$
assumes $0 < r$
shows $\exists b > 0. b \leq r \wedge (\exists B. B = (\sum_{i \in \text{Basis}. b *_R i) \wedge (\forall y \in \text{cbox } (x - B) (x + B). y \in \text{cball } x r))$
 $\langle \text{proof} \rangle$

lemma *Pair1-in-Basis*: $i \in \text{Basis} \implies (i, 0) \in \text{Basis}$
and *Pair2-in-Basis*: $i \in \text{Basis} \implies (0, i) \in \text{Basis}$
 $\langle \text{proof} \rangle$

lemma *le-real-sqrt-sumsq'* [*simp*]: $y \leq \text{sqrt } (x * x + y * y)$
 $\langle \text{proof} \rangle$

lemma *cball-Pair-split-subset*: $\text{cball } (a, b) c \subseteq \text{cball } a c \times \text{cball } b c$
 $\langle \text{proof} \rangle$

lemma *cball-times-subset*: $\text{cball } a (c/2) \times \text{cball } b (c/2) \subseteq \text{cball } (a, b) c$
 $\langle \text{proof} \rangle$

lemma *eventually-bound-pairE*:
assumes $\text{isCont } f (t0, x0)$
obtains B **where**
 $B \geq 1$
 $\text{eventually } (\lambda e. \forall x \in \text{cball } t0 e \times \text{cball } x0 e. \text{norm } (f x) \leq B) (\text{at-right } 0)$
 $\langle \text{proof} \rangle$

lemma
eventually-in-cballs:
assumes $d > 0 \ c > 0$
shows $\text{eventually } (\lambda e. \text{cball } t0 (c * e) \times (\text{cball } x0 e) \subseteq \text{cball } (t0, x0) d) (\text{at-right } 0)$
 $\langle \text{proof} \rangle$

lemma *cball-eq-sing'*:
fixes $x :: 'a::\{\text{metric-space}, \text{perfect-space}\}$
shows $\text{cball } x e = \{y\} \iff e = 0 \wedge x = y$
 $\langle \text{proof} \rangle$

locale *ll-on-open = interval T for T +*
fixes $f::\text{real} \Rightarrow 'a::\{\text{banach}, \text{heine-borel}\} \Rightarrow 'a$ **and** X
assumes *local-lipschitz*: $\text{local-lipschitz } T X f$
assumes *cont*: $\bigwedge x. x \in X \implies \text{continuous-on } T (\lambda t. f t x)$
assumes *open-domain*[*intro!*, *simp*]: $\text{open } T \ \text{open } X$
begin

all flows on closed segments

definition *csols* **where**

$csols\ t0\ x0 = \{(x, t1). \{t0--t1\} \subseteq T \wedge x\ t0 = x0 \wedge (x\ solves\ ode\ f)\ \{t0--t1\}\ X\}$

the maximal existence interval

definition $existence\ interval\ t0\ x0 = (\bigcup (x, t1) \in csols\ t0\ x0 . \{t0--t1\})$

witness flow

definition $csol\ t0\ x0 = (SOME\ csol. \forall t \in existence\ interval\ t0\ x0. (csol\ t, t) \in csols\ t0\ x0)$

unique flow

definition $flow\ where\ flow\ t0\ x0 = (\lambda t. if\ t \in existence\ interval\ t0\ x0\ then\ csol\ t0\ x0\ t\ t\ else\ 0)$

end

locale $ll\ on\ open\ it =$

general?:— TODO: why is this qualification necessary? It seems only because of $ll\ on\ open\ it\ T\ f\ X$

$ll\ on\ open + \mathbf{fixes}\ t0::real$

— if possible, all development should be done with $t0$ as explicit parameter for initial time: then it can be instantiated with 0 for autonomous ODEs

context $ll\ on\ open\ begin$

sublocale $ll\ on\ open\ it\ where\ t0 = t0\ for\ t0\ \langle proof \rangle$

sublocale $continuous\ rhs\ T\ X\ f$
 $\langle proof \rangle$

end

context $ll\ on\ open\ it\ begin$

lemma $ll\ on\ open\ rev[intro, simp]: ll\ on\ open\ (preflect\ t0\ 'T)\ (\lambda t. -\ f\ (preflect\ t0\ t))\ X$
 $\langle proof \rangle$

lemma $eventually\ lipschitz:$

assumes $t0 \in T\ x0 \in X\ c > 0$

obtains $L\ where$

$eventually\ (\lambda u. \forall t' \in cball\ t0\ (c * u) \cap T.$

$L\ lipschitz\ on\ (cball\ x0\ u \cap X)\ (\lambda y. f\ t'\ y))\ (at\ right\ 0)$

$\langle proof \rangle$

lemmas $continuous\ on\ Times\ f = continuous$

lemmas $continuous\ on\ f = continuous\ rhs\ comp$

lemma

lipschitz-on-compact:
assumes compact K $K \subseteq T$
assumes compact Y $Y \subseteq X$
obtains L **where** $\bigwedge t. t \in K \implies L\text{-lipschitz-on } Y (f t)$
 $\langle \text{proof} \rangle$

lemma *csols-empty-iff*: $csols\ t0\ x0 = \{\}$ $\longleftrightarrow t0 \notin T \vee x0 \notin X$
 $\langle \text{proof} \rangle$

lemma *csols-notempty*: $t0 \in T \implies x0 \in X \implies csols\ t0\ x0 \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *existence-ivl-empty-iff[simp]*: $existence\text{-ivl}\ t0\ x0 = \{\}$ $\longleftrightarrow t0 \notin T \vee x0 \notin X$
 $\langle \text{proof} \rangle$

lemma *existence-ivl-empty1[simp]*: $t0 \notin T \implies existence\text{-ivl}\ t0\ x0 = \{\}$
and *existence-ivl-empty2[simp]*: $x0 \notin X \implies existence\text{-ivl}\ t0\ x0 = \{\}$
 $\langle \text{proof} \rangle$

lemma *flow-undefined*:
shows $t0 \notin T \implies flow\ t0\ x0 = (\lambda-. 0)$
 $x0 \notin X \implies flow\ t0\ x0 = (\lambda-. 0)$
 $\langle \text{proof} \rangle$

lemma (**in** *ll-on-open*) *flow-eq-in-existence-ivlI*:
assumes $\bigwedge u. x0 \in X \implies u \in existence\text{-ivl}\ t0\ x0 \longleftrightarrow g\ u \in existence\text{-ivl}\ s0\ x0$
assumes $\bigwedge u. x0 \in X \implies u \in existence\text{-ivl}\ t0\ x0 \implies flow\ t0\ x0\ u = flow\ s0\ x0\ (g\ u)$
shows $flow\ t0\ x0 = (\lambda t. flow\ s0\ x0\ (g\ t))$
 $\langle \text{proof} \rangle$

2.7.2 Global maximal flow with local Lipschitz

lemma *local-unique-solution*:
assumes *iv-defined*: $t0 \in T\ x0 \in X$
obtains $et\ ex\ B\ L$
where $et > 0\ 0 < ex\ cball\ t0\ et \subseteq T\ cball\ x0\ ex \subseteq X$
unique-on-cylinder $t0\ (cball\ t0\ et)\ x0\ ex\ f\ B\ L$
 $\langle \text{proof} \rangle$

lemma *mem-existence-ivl-iv-defined*:
assumes $t \in existence\text{-ivl}\ t0\ x0$
shows $t0 \in T\ x0 \in X$
 $\langle \text{proof} \rangle$

lemma *csol-mem-csols*:
assumes $t \in existence\text{-ivl}\ t0\ x0$

shows $(\text{csol } t0 \ x0 \ t, t) \in \text{csols } t0 \ x0$
<proof>

lemma *csol*:

assumes $t \in \text{existence-ivl } t0 \ x0$
shows $t \in T \ \{t0 \ -- \ t\} \subseteq T \ \text{csol } t0 \ x0 \ t \ t0 = x0 \ (\text{csol } t0 \ x0 \ t \ \text{solves-ode } f)$
 $\{t0 \ -- \ t\} \ X$
<proof>

lemma *existence-ivl-initial-time-iff[simp]*: $t0 \in \text{existence-ivl } t0 \ x0 \longleftrightarrow t0 \in T \wedge x0 \in X$
<proof>

lemma *existence-ivl-initial-time*: $t0 \in T \implies x0 \in X \implies t0 \in \text{existence-ivl } t0 \ x0$
<proof>

lemmas *mem-existence-ivl-subset* = *csol(1)*

lemma *existence-ivl-subset*:

$\text{existence-ivl } t0 \ x0 \subseteq T$
<proof>

lemma *is-interval-existence-ivl[intro, simp]*: *is-interval* (*existence-ivl* $t0 \ x0$)
<proof>

lemma *connected-existence-ivl[intro, simp]*: *connected* (*existence-ivl* $t0 \ x0$)
<proof>

lemma *in-existence-between-zeroI*:

$t \in \text{existence-ivl } t0 \ x0 \implies s \in \{t0 \ -- \ t\} \implies s \in \text{existence-ivl } t0 \ x0$
<proof>

lemma *segment-subset-existence-ivl*:

assumes $s \in \text{existence-ivl } t0 \ x0 \ t \in \text{existence-ivl } t0 \ x0$
shows $\{s \ -- \ t\} \subseteq \text{existence-ivl } t0 \ x0$
<proof>

lemma *flow-initial-time-if*: $\text{flow } t0 \ x0 \ t0 = (\text{if } t0 \in T \wedge x0 \in X \text{ then } x0 \text{ else } 0)$
<proof>

lemma *flow-initial-time[simp]*: $t0 \in T \implies x0 \in X \implies \text{flow } t0 \ x0 \ t0 = x0$
<proof>

lemma *open-existence-ivl[intro, simp]*: *open* (*existence-ivl* $t0 \ x0$)
<proof>

lemma *csols-unique*:

assumes $(x, t1) \in \text{csols } t0 \ x0$
assumes $(y, t2) \in \text{csols } t0 \ x0$

shows $\forall t \in \{t0 \text{ -- } t1\} \cap \{t0 \text{ -- } t2\}. x t = y t$
<proof>

lemma *csol-unique:*

assumes *t1:* $t1 \in \text{existence-ivl } t0 \ x0$
assumes *t2:* $t2 \in \text{existence-ivl } t0 \ x0$
assumes *t:* $t \in \{t0 \text{ -- } t1\} \ t \in \{t0 \text{ -- } t2\}$
shows $\text{csol } t0 \ x0 \ t1 \ t = \text{csol } t0 \ x0 \ t2 \ t$
<proof>

lemma *flow-vderiv-on-left:*

$(\text{flow } t0 \ x0 \ \text{has-vderiv-on } (\lambda x. f \ x \ (\text{flow } t0 \ x0 \ x))) \ (\text{existence-ivl } t0 \ x0 \cap \{..t0\})$
<proof>

lemma *flow-vderiv-on-right:*

$(\text{flow } t0 \ x0 \ \text{has-vderiv-on } (\lambda x. f \ x \ (\text{flow } t0 \ x0 \ x))) \ (\text{existence-ivl } t0 \ x0 \cap \{t0..\})$
<proof>

lemma *flow-usolves-ode:*

assumes *iv-defined:* $t0 \in T \ x0 \in X$
shows $(\text{flow } t0 \ x0 \ \text{usolves-ode } f \ \text{from } t0) \ (\text{existence-ivl } t0 \ x0) \ X$
<proof>

lemma *flow-solves-ode:* $t0 \in T \implies x0 \in X \implies (\text{flow } t0 \ x0 \ \text{solves-ode } f) \ (\text{existence-ivl } t0 \ x0) \ X$
<proof>

lemma *equals-flowI:*

assumes $t0 \in T'$
is-interval T'
 $T' \subseteq \text{existence-ivl } t0 \ x0$
 $(z \ \text{solves-ode } f) \ T' \ X$
 $z \ t0 = \text{flow } t0 \ x0 \ t0 \ t \in T'$
shows $z \ t = \text{flow } t0 \ x0 \ t$
<proof>

lemma *existence-ivl-maximal-segment:*

assumes $(x \ \text{solves-ode } f) \ \{t0 \text{ -- } t\} \ X \ x \ t0 = x0$
assumes $\{t0 \text{ -- } t\} \subseteq T$
shows $t \in \text{existence-ivl } t0 \ x0$
<proof>

lemma *existence-ivl-maximal-interval:*

assumes $(x \ \text{solves-ode } f) \ S \ X \ x \ t0 = x0$
assumes $t0 \in S$ *is-interval* $S \ S \subseteq T$
shows $S \subseteq \text{existence-ivl } t0 \ x0$
<proof>

lemma *maximal-existence-flow:*

assumes *sol*: (*x solves-ode f*) *K X* **and** *iv*: $x\ t0 = x0$
assumes *is-interval K*
assumes $t0 \in K$
assumes $K \subseteq T$
shows $K \subseteq \text{existence-ivl } t0\ x0 \wedge t. t \in K \implies \text{flow } t0\ x0\ t = x\ t$
 <proof>

lemma *maximal-existence-flowI*:
assumes (*x has-vderiv-on* ($\lambda t. f\ t\ (x\ t)$)) *K*
assumes $\wedge t. t \in K \implies x\ t \in X$
assumes $x\ t0 = x0$
assumes *K*: *is-interval K* $t0 \in K$ $K \subseteq T$
shows $K \subseteq \text{existence-ivl } t0\ x0 \wedge t. t \in K \implies \text{flow } t0\ x0\ t = x\ t$
 <proof>

lemma *flow-in-domain*: $t \in \text{existence-ivl } t0\ x0 \implies \text{flow } t0\ x0\ t \in X$
 <proof>

lemma (*in ll-on-open*)
assumes $t \in \text{existence-ivl } s\ x$
assumes $x \in X$
assumes *auto*: $\wedge s\ t\ x. x \in X \implies f\ s\ x = f\ t\ x$
assumes $T = \text{UNIV}$
shows *mem-existence-ivl-shift-autonomous1*: $t - s \in \text{existence-ivl } 0\ x$
and *flow-shift-autonomous1*: $\text{flow } s\ x\ t = \text{flow } 0\ x\ (t - s)$
 <proof>

lemma (*in ll-on-open*)
assumes $t - s \in \text{existence-ivl } 0\ x$
assumes $x \in X$
assumes *auto*: $\wedge s\ t\ x. x \in X \implies f\ s\ x = f\ t\ x$
assumes $T = \text{UNIV}$
shows *mem-existence-ivl-shift-autonomous2*: $t \in \text{existence-ivl } s\ x$
and *flow-shift-autonomous2*: $\text{flow } s\ x\ t = \text{flow } 0\ x\ (t - s)$
 <proof>

lemma
flow-eq-rev:
assumes $t \in \text{existence-ivl } t0\ x0$
shows *preflect* $t0\ t \in \text{ll-on-open.existence-ivl } (\text{preflect } t0\ 'T) (\lambda t. - f (\text{preflect } t0\ t))\ X\ t0\ x0$
 $\text{flow } t0\ x0\ t = \text{ll-on-open.flow } (\text{preflect } t0\ 'T) (\lambda t. - f (\text{preflect } t0\ t))\ X\ t0\ x0$
 (<proof> *preflect* $t0\ t$)
 <proof>

lemma (*in ll-on-open*)
shows *rev-flow-eq*: $t \in \text{ll-on-open.existence-ivl } (\text{preflect } t0\ 'T) (\lambda t. - f (\text{preflect } t0\ t))\ X\ t0\ x0 \implies$
 $\text{ll-on-open.flow } (\text{preflect } t0\ 'T) (\lambda t. - f (\text{preflect } t0\ t))\ X\ t0\ x0\ t = \text{flow } t0\ x0\ t$

(*preflect t0 t*)
and *mem-rev-existence-ivl-eq*:
 $t \in ll\text{-on-open.existence-ivl } (preflect\ t0\ 'T) (\lambda t. - f (preflect\ t0\ t)) X\ t0\ x0 \longleftrightarrow$
 $preflect\ t0\ t \in existence\text{-ivl } t0\ x0$
<proof>

lemma

shows *rev-existence-ivl-eq*: $ll\text{-on-open.existence-ivl } (preflect\ t0\ 'T) (\lambda t. - f (preflect\ t0\ t)) X\ t0\ x0 = preflect\ t0\ 'existence\text{-ivl } t0\ x0$
and *existence-ivl-eq-rev*: $existence\text{-ivl } t0\ x0 = preflect\ t0\ 'll\text{-on-open.existence-ivl } (preflect\ t0\ 'T) (\lambda t. - f (preflect\ t0\ t)) X\ t0\ x0$
<proof>

end

end

3 Bounded Linear Operator

theory *Bounded-Linear-Operator*

imports

HOL-Analysis.Analysis

begin

typedef (**overloaded**) *'a blinop* = *UNIV::('a, 'a) blinfun set*
<proof>

setup-lifting *type-definition-blinop*

lift-definition *blinop-apply::('a::real-normed-vector) blinop \Rightarrow 'a \Rightarrow 'a* **is** *blinfun-apply* *<proof>*

lift-definition *Blinop::('a::real-normed-vector \Rightarrow 'a) \Rightarrow 'a blinop* **is** *Blinfun* *<proof>*

no-notation *vec-nth* (**infixl** $\langle \$ \rangle$ 90)

notation *blinop-apply* (**infixl** $\langle \$ \rangle$ 999)

declare $[[coercion\ blinop\text{-}apply :: ('a::real\text{-}normed\text{-}vector)\ blinop \Rightarrow 'a \Rightarrow 'a]]$

instantiation *blinop :: (real-normed-vector) real-normed-vector*

begin

lift-definition *norm-blinop :: 'a blinop \Rightarrow real* **is** *norm* *<proof>*

lift-definition *minus-blinop :: 'a blinop \Rightarrow 'a blinop \Rightarrow 'a blinop* **is** *minus* *<proof>*

lift-definition *dist-blinop :: 'a blinop \Rightarrow 'a blinop \Rightarrow real* **is** *dist* *<proof>*

definition *uniformity-blinop :: ('a blinop \times 'a blinop) filter* **where**

uniformity-blinop = (INF e \in {0<..}. principal {(x, y). dist x y < e})

definition *open-blinop* :: 'a blinop set \Rightarrow bool **where**
open-blinop U = ($\forall x \in U. \forall_F (x', y)$ in uniformity. $x' = x \longrightarrow y \in U$)

lift-definition *uminus-blinop* :: 'a blinop \Rightarrow 'a blinop **is** *uminus* \langle proof \rangle

lift-definition *zero-blinop* :: 'a blinop **is** 0 \langle proof \rangle

lift-definition *plus-blinop* :: 'a blinop \Rightarrow 'a blinop \Rightarrow 'a blinop **is** *plus* \langle proof \rangle

lift-definition *scaleR-blinop*::real \Rightarrow 'a blinop \Rightarrow 'a blinop **is** *scaleR* \langle proof \rangle

lift-definition *sgn-blinop* :: 'a blinop \Rightarrow 'a blinop **is** *sgn* \langle proof \rangle

instance
 \langle proof \rangle
end

lemma *bounded-bilinear-blinop-apply*: bounded-bilinear (\$) \langle proof \rangle

interpretation *blinop*: bounded-bilinear (\$) \langle proof \rangle

lemma *blinop-eqI*: ($\bigwedge i. x \$ i = y \$ i$) \Longrightarrow $x = y$
 \langle proof \rangle

lemmas *bounded-linear-apply-blinop*[intro, simp] = *blinop.bounded-linear-left*
declare *blinop.tendsto*[tendsto-intros]
declare *blinop.FDERIV*[derivative-intros]
declare *blinop.continuous*[continuous-intros]
declare *blinop.continuous-on*[continuous-intros]

instance *blinop* :: (banach) banach
 \langle proof \rangle

instance *blinop* :: (euclidean-space) heine-borel
 \langle proof \rangle

instantiation *blinop*::({real-normed-vector, perfect-space}) real-normed-algebra-1
begin

lift-definition *one-blinop*::'a blinop **is** *id-blinfun* \langle proof \rangle

lemma *blinop-apply-one-blinop*[simp]: 1 \$ x = x
 \langle proof \rangle

lift-definition *times-blinop* :: 'a blinop \Rightarrow 'a blinop \Rightarrow 'a blinop **is** *blinfun-compose*
 \langle proof \rangle

lemma *blinop-apply-times-blinop*[simp]: $(f * g) \$ x = f \$ (g \$ x)$
 ⟨proof⟩

instance
 ⟨proof⟩
end

lemmas *bounded-bilinear-bounded-uniform-limit-intros*[*uniform-limit-intros*] =
bounded-bilinear.bounded-uniform-limit[*OF Bounded-Linear-Operator.bounded-bilinear-blinop-apply*]
bounded-bilinear.bounded-uniform-limit[*OF Bounded-Linear-Function.bounded-bilinear-blinfun-apply*]
bounded-bilinear.bounded-uniform-limit[*OF Bounded-Linear-Operator.blinop.flip*]
bounded-bilinear.bounded-uniform-limit[*OF Bounded-Linear-Function.blinfun.flip*]
bounded-linear.uniform-limit[*OF blinop.bounded-linear-right*]
bounded-linear.uniform-limit[*OF blinop.bounded-linear-left*]
bounded-linear.uniform-limit[*OF bounded-linear-apply-blinop*]

no-notation
blinop-apply (**infixl** <\$> 999)
notation *vec-nth* (**infixl** <\$> 90)

end

4 Multivariate Taylor

theory *Multivariate-Taylor*
imports
HOL-Analysis.Analysis
 ../*ODE-Auxiliarities*
begin

no-notation *vec-nth* (**infixl** <\$> 90)
notation *blinfun-apply* (**infixl** <\$> 999)

lemma
fixes $f::'a::\text{real-normed-vector} \Rightarrow 'b::\text{banach}$
and $Df::'a \Rightarrow \text{nat} \Rightarrow 'a \Rightarrow 'a \Rightarrow 'b$
assumes $n > 0$
assumes *Df-Nil*: $\bigwedge a x. Df a 0 H H = f a$
assumes *Df-Cons*: $\bigwedge a i d. a \in \text{closed-segment } X (X + H) \implies i < n \implies$
 $((\lambda a. Df a i H H) \text{ has-derivative } (Df a (Suc i) H)) \text{ (at } a \text{ within } G)$
assumes *cs*: $\text{closed-segment } X (X + H) \subseteq G$
defines $i \equiv \lambda x.$
 $((1 - x) ^ (n - 1) / \text{fact } (n - 1)) *_R Df (X + x *_R H) n H H$
shows *multivariate-Taylor-has-integral*:
 $(i \text{ has-integral } f (X + H) - (\sum i < n. (1 / \text{fact } i) *_R Df X i H H)) \{0..1\}$
and *multivariate-Taylor*:
 $f (X + H) = (\sum i < n. (1 / \text{fact } i) *_R Df X i H H) + \text{integral } \{0..1\} i$
and *multivariate-Taylor-integrable*:
 $i \text{ integrable-on } \{0..1\}$

<proof>

4.1 Symmetric second derivative

lemma *symmetric-second-derivative-aux*:
 assumes *first-fderiv[derivative-intros]*:
 $\bigwedge a. a \in G \implies (f \text{ has-derivative } (f' a)) \text{ (at } a \text{ within } G)$
 assumes *second-fderiv[derivative-intros]*:
 $\bigwedge i. ((\lambda x. f' x i) \text{ has-derivative } (\lambda j. f'' j i)) \text{ (at } a \text{ within } G)$
 assumes $i \neq j \ i \neq 0 \ j \neq 0$
 assumes $a \in G$
 assumes $\bigwedge s \ t. s \in \{0..1\} \implies t \in \{0..1\} \implies a + s *_R i + t *_R j \in G$
 shows $f'' j i = f'' i j$
<proof>

locale *second-derivative-within* =
 fixes $f \ f' \ f'' \ a \ G$
 assumes *first-fderiv[derivative-intros]*:
 $\bigwedge a. a \in G \implies (f \text{ has-derivative } \text{blinfun-apply } (f' a)) \text{ (at } a \text{ within } G)$
 assumes *in-G*: $a \in G$
 assumes *second-fderiv[derivative-intros]*:
 $(f' \text{ has-derivative } \text{blinfun-apply } f'')$ (at a within G)
begin

lemma *symmetric-second-derivative-within*:
 assumes $a \in G$
 assumes $\bigwedge s \ t. s \in \{0..1\} \implies t \in \{0..1\} \implies a + s *_R i + t *_R j \in G$
 shows $f'' i j = f'' j i$
<proof>

end

locale *second-derivative* =
 fixes $f :: 'a :: \text{real-normed-vector} \Rightarrow 'b :: \text{banach}$
 and $f' :: 'a \Rightarrow 'a \Rightarrow_L 'b$
 and $f'' :: 'a \Rightarrow_L 'a \Rightarrow_L 'b$
 and $a :: 'a$
 and $G :: 'a \text{ set}$
 assumes *first-fderiv[derivative-intros]*:
 $\bigwedge a. a \in G \implies (f \text{ has-derivative } f' a) \text{ (at } a)$
 assumes *in-G*: $a \in \text{interior } G$
 assumes *second-fderiv[derivative-intros]*:
 $(f' \text{ has-derivative } f'')$ (at a)
begin

lemma *symmetric-second-derivative*:
 assumes $a \in \text{interior } G$
 shows $f'' i j = f'' j i$
<proof>

end

lemma

uniform-explicit-remainder-Taylor-1:

fixes $f::'a::\{\text{banach,heine-borel,perfect-space}\} \Rightarrow 'b::\text{banach}$

assumes f' [*derivative-intros*]: $\bigwedge x. x \in G \Longrightarrow (f \text{ has-derivative } \text{blinfun-apply } (f' x)) \text{ (at } x)$

assumes f' -*cont*: $\bigwedge x. x \in G \Longrightarrow \text{isCont } f' x$

assumes *open* G

assumes $J \neq \{\}$ *compact* J $J \subseteq G$

assumes $e > 0$

obtains $d R$

where $d > 0$

$\bigwedge x z. f z = f x + f' x (z - x) + R x z$

$\bigwedge x y. x \in J \Longrightarrow y \in J \Longrightarrow \text{dist } x y < d \Longrightarrow \text{norm } (R x y) \leq e * \text{dist } x y$

continuous-on $(G \times G) (\lambda(a, b). R a b)$

<proof>

TODO: rename, duplication?

locale *second-derivative-within'* =

fixes $f f' f'' a G$

assumes f [*first-fderiv*][*derivative-intros*]:

$\bigwedge a. a \in G \Longrightarrow (f \text{ has-derivative } f' a) \text{ (at } a \text{ within } G)$

assumes *in-G*: $a \in G$

assumes f [*second-fderiv*][*derivative-intros*]:

$\bigwedge i. ((\lambda x. f' x i) \text{ has-derivative } f'' i) \text{ (at } a \text{ within } G)$

begin

lemma *symmetric-second-derivative-within*:

assumes $a \in G$ *open* G

assumes $\bigwedge s t. s \in \{0..1\} \Longrightarrow t \in \{0..1\} \Longrightarrow a + s *_R i + t *_R j \in G$

shows $f'' i j = f'' j i$

<proof>

end

locale *second-derivative-on-open* =

fixes $f::'a::\text{real-normed-vector} \Rightarrow 'b::\text{banach}$

and $f'::'a \Rightarrow 'a \Rightarrow 'b$

and $f''::'a \Rightarrow 'a \Rightarrow 'b$

and $a::'a$

and $G::'a \text{ set}$

assumes f [*first-fderiv*][*derivative-intros*]:

$\bigwedge a. a \in G \Longrightarrow (f \text{ has-derivative } f' a) \text{ (at } a)$

assumes *in-G*: $a \in G$ **and** *open-G*: *open* G

assumes f [*second-fderiv*][*derivative-intros*]:

$((\lambda x. f' x i) \text{ has-derivative } f'' i) \text{ (at } a)$

begin

lemma *symmetric-second-derivative*:

assumes $a \in G$

shows $f''\ i\ j = f''\ j\ i$

<proof>

end

no-notation

blinfun-apply (**infixl** $\langle \$ \rangle$ 999)

notation *vec-nth* (**infixl** $\langle \$ \rangle$ 90)

end

5 Flow

theory *Flow*

imports

Picard-Lindelof-Qualitative

HOL-Library.Diagonal-Subsequence

../Library/Bounded-Linear-Operator

../Library/Multivariate-Taylor

../Library/Interval-Integral-HK

begin

TODO: extend theorems for dependence on initial time

5.1 simp rules for integrability (TODO: move)

lemma *blinfun-ext*: $x = y \longleftrightarrow (\forall i. \text{blinfun-apply } x\ i = \text{blinfun-apply } y\ i)$

<proof>

notation *id-blinfun* ($\langle 1_L \rangle$)

lemma *blinfun-inverse-left*:

fixes $f::'a::\text{euclidean-space} \Rightarrow_L 'a$ **and** f'

shows $f\ o_L\ f' = 1_L \longleftrightarrow f'\ o_L\ f = 1_L$

<proof>

lemma *onorm-zero-blinfun[simp]*: $\text{onorm } (\text{blinfun-apply } 0) = 0$

<proof>

lemma *blinfun-compose-1-left[simp]*: $x\ o_L\ 1_L = x$

and *blinfun-compose-1-right[simp]*: $1_L\ o_L\ y = y$

<proof>

named-theorems *integrable-on-simps*

lemma *integrable-on-refl-ivl*[*intro, simp*]: g *integrable-on* $\{b .. (b::'b::ordered-euclidean-space)\}$
and *integrable-on-refl-closed-segment*[*intro, simp*]: h *integrable-on closed-segment*
a a
<proof>

lemma *integrable-const-ivl-closed-segment*[*intro, simp*]: $(\lambda x. c)$ *integrable-on closed-segment*
a (b::real)
<proof>

lemma *integrable-ident-ivl*[*intro, simp*]: $(\lambda x. x)$ *integrable-on closed-segment a (b::real)*
and *integrable-ident-cbox*[*intro, simp*]: $(\lambda x. x)$ *integrable-on cbox a (b::real)*
<proof>

lemma *content-closed-segment-real*:
fixes *a b::real*
shows *content (closed-segment a b) = abs (b - a)*
<proof>

lemma *integral-const-closed-segment*:
fixes *a b::real*
shows *integral (closed-segment a b) (\lambda x. c) = abs (b - a) *_R c*
<proof>

lemmas [*integrable-on-simps*] =
integrable-on-empty — *empty*
integrable-on-refl integrable-on-refl-ivl integrable-on-refl-closed-segment — *singleton*
integrable-const integrable-const-ivl integrable-const-ivl-closed-segment — *constant*
ident-integrable-on integrable-ident-ivl integrable-ident-cbox — *identity*

lemma *integrable-cmul-real*:
fixes *K::real*
shows *f integrable-on X \implies (\lambda x. K * f x) integrable-on X*
<proof>

lemmas [*integrable-on-simps*] =
integrable-0
integrable-neg
integrable-cmul
integrable-cmul-real
integrable-on-cmult-iff
integrable-on-cmult-left
integrable-on-cmult-right
integrable-on-cmult-iff
integrable-on-cmult-left-iff
integrable-on-cmult-right-iff
integrable-on-cdivide-iff
integrable-diff
integrable-add

integrable-sum

lemma *dist-cancel-add1*: $\text{dist } (t0 + et) t0 = \text{norm } et$
<proof>

lemma *double-nonneg-le*:
fixes $a::\text{real}$
shows $a * 2 \leq b \implies a \geq 0 \implies a \leq b$
<proof>

5.2 Nonautonomous IVP on maximal existence interval

context *ll-on-open-it*
begin

context
fixes $x0$
assumes *iv-defined*: $t0 \in T \ x0 \in X$
begin

lemmas *closed-segment-iv-subset-domain* = *closed-segment-subset-domainI*[*OF iv-defined*(1)]

lemma
local-unique-solutions:
obtains $t \ u \ L$
where
 $0 < t0 < u$
 $\text{cball } t0 \ t \subseteq \text{existence-ivl } t0 \ x0$
 $\text{cball } x0 \ (2 * u) \subseteq X$
 $\bigwedge t'. t' \in \text{cball } t0 \ t \implies L\text{-lipschitz-on } (\text{cball } x0 \ (2 * u)) \ (f \ t')$
 $\bigwedge x. x \in \text{cball } x0 \ u \implies (\text{flow } t0 \ x \ \text{usolves-ode } f \ \text{from } t0) \ (\text{cball } t0 \ t) \ (\text{cball } x \ u)$
 $\bigwedge x. x \in \text{cball } x0 \ u \implies \text{cball } x \ u \subseteq X$
<proof>

lemma *Picard-iterate-mem-existence-ivlI*:
assumes $t \in T$
assumes *compact* $C \ x0 \in C \ C \subseteq X$
assumes $\bigwedge y \ s. s \in \{t0 \ \text{--} \ t\} \implies y \ t0 = x0 \implies y \in \{t0 \ \text{--} \ s\} \rightarrow C \implies$
continuous-on $\{t0 \ \text{--} \ s\} \ y \implies$
 $x0 + \text{ivl-integral } t0 \ s \ (\lambda t. f \ t \ (y \ t)) \in C$
shows $t \in \text{existence-ivl } t0 \ x0 \ \bigwedge s. s \in \{t0 \ \text{--} \ t\} \implies \text{flow } t0 \ x0 \ s \in C$
<proof>

lemma *flow-has-vderiv-on*: $(\text{flow } t0 \ x0 \ \text{has-vderiv-on } (\lambda t. f \ t \ (\text{flow } t0 \ x0 \ t))) \ (\text{existence-ivl } t0 \ x0)$
<proof>

lemmas *flow-has-vderiv-on-compose*[*derivative-intros*] =
has-vderiv-on-compose2[*OF flow-has-vderiv-on, THEN has-vderiv-on-eq-rhs*]

end

lemma *unique-on-intersection*:

assumes *sols*: $(x \text{ solves-ode } f) U X (y \text{ solves-ode } f) V X$
assumes *iv-mem*: $t0 \in U t0 \in V$ **and** *subs*: $U \subseteq T V \subseteq T$
assumes *ivls*: *is-interval* U *is-interval* V
assumes *iv*: $x t0 = y t0$
assumes *mem*: $t \in U t \in V$
shows $x t = y t$

<proof>

lemma *unique-solution*:

assumes *sols*: $(x \text{ solves-ode } f) U X (y \text{ solves-ode } f) U X$
assumes *iv-mem*: $t0 \in U$ **and** *subs*: $U \subseteq T$
assumes *ivls*: *is-interval* U
assumes *iv*: $x t0 = y t0$
assumes *mem*: $t \in U$
shows $x t = y t$

<proof>

lemma

assumes *s*: $s \in \text{existence-ivl } t0 x0$
assumes *t*: $t + s \in \text{existence-ivl } s (\text{flow } t0 x0 s)$
shows *flow-trans*: $\text{flow } t0 x0 (s + t) = \text{flow } s (\text{flow } t0 x0 s) (s + t)$
and *existence-ivl-trans*: $s + t \in \text{existence-ivl } t0 x0$

<proof>

lemma

assumes *t*: $t \in \text{existence-ivl } t0 x0$
shows *flows-reverse*: $\text{flow } t (\text{flow } t0 x0 t) t0 = x0$
and *existence-ivl-reverse*: $t0 \in \text{existence-ivl } t (\text{flow } t0 x0 t)$

<proof>

lemma *flow-has-derivative*:

assumes *t* $\in \text{existence-ivl } t0 x0$
shows $(\text{flow } t0 x0 \text{ has-derivative } (\lambda i. i *_R f t (\text{flow } t0 x0 t))) (at t)$

<proof>

lemma *flow-has-vector-derivative*:

assumes *t* $\in \text{existence-ivl } t0 x0$
shows $(\text{flow } t0 x0 \text{ has-vector-derivative } f t (\text{flow } t0 x0 t)) (at t)$

<proof>

lemma *flow-has-vector-derivative-at-0*:

assumes *t* $\in \text{existence-ivl } t0 x0$
shows $((\lambda h. \text{flow } t0 x0 (t + h)) \text{ has-vector-derivative } f t (\text{flow } t0 x0 t)) (at 0)$

<proof>

lemma

assumes $t \in \text{existence-ivl } t0 \ x0$

shows *closed-segment-subset-existence-ivl*: $\text{closed-segment } t0 \ t \subseteq \text{existence-ivl } t0 \ x0$

and *ivl-subset-existence-ivl*: $\{t0 \ .. \ t\} \subseteq \text{existence-ivl } t0 \ x0$

and *ivl-subset-existence-ivl'*: $\{t \ .. \ t0\} \subseteq \text{existence-ivl } t0 \ x0$

<proof>

lemma *flow-fixed-point*:

assumes $t: t \in \text{existence-ivl } t0 \ x0$

shows $\text{flow } t0 \ x0 \ t = x0 + \text{ivl-integral } t0 \ t \ (\lambda t. f \ t \ (\text{flow } t0 \ x0 \ t))$

<proof>

lemma *flow-continuous*: $t \in \text{existence-ivl } t0 \ x0 \implies \text{continuous } (\text{at } t) \ (\text{flow } t0 \ x0)$

<proof>

lemma *flow-tendsto*: $t \in \text{existence-ivl } t0 \ x0 \implies (ts \longrightarrow t) \ F \implies$

$((\lambda s. \text{flow } t0 \ x0 \ (ts \ s)) \longrightarrow \text{flow } t0 \ x0 \ t) \ F$

<proof>

lemma *flow-continuous-on*: $\text{continuous-on } (\text{existence-ivl } t0 \ x0) \ (\text{flow } t0 \ x0)$

<proof>

lemma *flow-continuous-on-intro*:

continuous-on $s \ g \implies$

$(\bigwedge xa. xa \in s \implies g \ xa \in \text{existence-ivl } t0 \ x0) \implies$

continuous-on $s \ (\lambda xa. \text{flow } t0 \ x0 \ (g \ xa))$

<proof>

lemma *f-flow-continuous*:

assumes $t \in \text{existence-ivl } t0 \ x0$

shows $\text{isCont } (\lambda t. f \ t \ (\text{flow } t0 \ x0 \ t)) \ t$

<proof>

lemma *exponential-initial-condition*:

assumes $y0: t \in \text{existence-ivl } t0 \ y0$

assumes $z0: t \in \text{existence-ivl } t0 \ z0$

assumes $Y \subseteq X$

assumes *remain*: $\bigwedge s. s \in \text{closed-segment } t0 \ t \implies \text{flow } t0 \ y0 \ s \in Y$

$\bigwedge s. s \in \text{closed-segment } t0 \ t \implies \text{flow } t0 \ z0 \ s \in Y$

assumes *lipschitz*: $\bigwedge s. s \in \text{closed-segment } t0 \ t \implies K\text{-lipschitz-on } Y \ (f \ s)$

shows $\text{norm } (\text{flow } t0 \ y0 \ t - \text{flow } t0 \ z0 \ t) \leq \text{norm } (y0 - z0) * \exp ((K + 1) * \text{abs } (t - t0))$

<proof>

lemma

existence-ivl-cballs:

assumes *iv-defined*: $t0 \in T \ x0 \in X$

obtains $t \ u \ L$

where

$\bigwedge y. y \in \text{cball } x0 \ u \implies \text{cball } t0 \ t \subseteq \text{existence-ivl } t0 \ y$

$\bigwedge s \ y. y \in \text{cball } x0 \ u \implies s \in \text{cball } t0 \ t \implies \text{flow } t0 \ y \ s \in \text{cball } y \ u$

$L\text{-lipschitz-on } (\text{cball } t0 \ t \times \text{cball } x0 \ u) \ (\lambda(t, x). \text{flow } t0 \ x \ t)$

$\bigwedge y. y \in \text{cball } x0 \ u \implies \text{cball } y \ u \subseteq X$

$0 < t0 < u$

$\langle \text{proof} \rangle$

context

fixes $x0$

assumes $\text{iv-defined}: t0 \in T \ x0 \in X$

begin

lemma $\text{existence-ivl-notempty}: \text{existence-ivl } t0 \ x0 \neq \{\}$

$\langle \text{proof} \rangle$

lemma $\text{initial-time-bounds}$:

shows $\text{bdd-above } (\text{existence-ivl } t0 \ x0) \implies t0 < \text{Sup } (\text{existence-ivl } t0 \ x0)$ (**is** $?a \implies -$)

and $\text{bdd-below } (\text{existence-ivl } t0 \ x0) \implies \text{Inf } (\text{existence-ivl } t0 \ x0) < t0$ (**is** $?b \implies -$)

$\langle \text{proof} \rangle$

lemma

$\text{flow-leaves-compact-ivl-right}$:

assumes $\text{bdd}: \text{bdd-above } (\text{existence-ivl } t0 \ x0)$

defines $b \equiv \text{Sup } (\text{existence-ivl } t0 \ x0)$

assumes $b \in T$

assumes $\text{compact } K$

assumes $K \subseteq X$

obtains t **where** $t \geq t0 \ t \in \text{existence-ivl } t0 \ x0 \ \text{flow } t0 \ x0 \ t \notin K$

$\langle \text{proof} \rangle$

lemma

$\text{flow-leaves-compact-ivl-left}$:

assumes $\text{bdd}: \text{bdd-below } (\text{existence-ivl } t0 \ x0)$

defines $b \equiv \text{Inf } (\text{existence-ivl } t0 \ x0)$

assumes $b \in T$

assumes $\text{compact } K$

assumes $K \subseteq X$

obtains t **where** $t \leq t0 \ t \in \text{existence-ivl } t0 \ x0 \ \text{flow } t0 \ x0 \ t \notin K$

$\langle \text{proof} \rangle$

lemma

$\text{sup-existence-maximal}$:

assumes $\bigwedge t. t0 \leq t \implies t \in \text{existence-ivl } t0 \ x0 \implies \text{flow } t0 \ x0 \ t \in K$

assumes $\text{compact } K \ K \subseteq X$

assumes $\text{bdd-above } (\text{existence-ivl } t0 \ x0)$

shows $Sup (existence-ivl\ t0\ x0) \notin T$
<proof>

lemma

inf-existence-minimal:

assumes $\bigwedge t. t \leq t0 \implies t \in existence-ivl\ t0\ x0 \implies flow\ t0\ x0\ t \in K$

assumes *compact* $K\ K \subseteq X$

assumes *bdd-below* $(existence-ivl\ t0\ x0)$

shows $Inf (existence-ivl\ t0\ x0) \notin T$

<proof>

end

lemma

subset-mem-compact-implies-subset-existence-interval:

assumes *ivl*: $t0 \in T'$ *is-interval* $T'\ T' \subseteq T$

assumes *iv-defined*: $x0 \in X$

assumes *mem-compact*: $\bigwedge t. t \in T' \implies t \in existence-ivl\ t0\ x0 \implies flow\ t0\ x0\ t \in K$

assumes *K*: *compact* $K\ K \subseteq X$

shows $T' \subseteq existence-ivl\ t0\ x0$

<proof>

lemma

mem-compact-implies-subset-existence-interval:

assumes *iv-defined*: $t0 \in T\ x0 \in X$

assumes *mem-compact*: $\bigwedge t. t \in T \implies t \in existence-ivl\ t0\ x0 \implies flow\ t0\ x0\ t \in K$

assumes *K*: *compact* $K\ K \subseteq X$

shows $T \subseteq existence-ivl\ t0\ x0$

<proof>

lemma

global-right-existence-ivl-explicit:

assumes $b \geq t0$

assumes *b*: $b \in existence-ivl\ t0\ x0$

obtains $d\ K$ **where** $d > 0\ K > 0$

$ball\ x0\ d \subseteq X$

$\bigwedge y. y \in ball\ x0\ d \implies b \in existence-ivl\ t0\ y$

$\bigwedge t\ y. y \in ball\ x0\ d \implies t \in \{t0 .. b\} \implies$

$dist (flow\ t0\ x0\ t) (flow\ t0\ y\ t) \leq dist\ x0\ y * exp (K * abs (t - t0))$

<proof>

lemma

global-left-existence-ivl-explicit:

assumes $b \leq t0$

assumes *b*: $b \in existence-ivl\ t0\ x0$

assumes *iv-defined*: $t0 \in T\ x0 \in X$

obtains $d\ K$ **where** $d > 0\ K > 0$

$ball\ x0\ d \subseteq X$
 $\bigwedge y. y \in ball\ x0\ d \implies b \in existence\text{-}ivl\ t0\ y$
 $\bigwedge t\ y. y \in ball\ x0\ d \implies t \in \{b .. t0\} \implies dist\ (flow\ t0\ x0\ t)\ (flow\ t0\ y\ t) \leq dist\ x0\ y * exp\ (K * abs\ (t - t0))$
 <proof>

lemma

global-existence-ivl-explicit:

assumes $a: a \in existence\text{-}ivl\ t0\ x0$

assumes $b: b \in existence\text{-}ivl\ t0\ x0$

assumes $le: a \leq b$

obtains $d\ K$ **where** $d > 0\ K > 0$

$ball\ x0\ d \subseteq X$

$\bigwedge y. y \in ball\ x0\ d \implies a \in existence\text{-}ivl\ t0\ y$

$\bigwedge y. y \in ball\ x0\ d \implies b \in existence\text{-}ivl\ t0\ y$

$\bigwedge t\ y. y \in ball\ x0\ d \implies t \in \{a .. b\} \implies$

$dist\ (flow\ t0\ x0\ t)\ (flow\ t0\ y\ t) \leq dist\ x0\ y * exp\ (K * abs\ (t - t0))$

<proof>

lemma *eventually-exponential-separation:*

assumes $a: a \in existence\text{-}ivl\ t0\ x0$

assumes $b: b \in existence\text{-}ivl\ t0\ x0$

assumes $le: a \leq b$

obtains K **where** $K > 0\ \forall_F\ y\ in\ at\ x0. \forall t \in \{a..b\}. dist\ (flow\ t0\ x0\ t)\ (flow\ t0\ y\ t) \leq dist\ x0\ y * exp\ (K * |t - t0|)$

<proof>

lemma *eventually-mem-existence-ivl:*

assumes $b: b \in existence\text{-}ivl\ t0\ x0$

shows $\forall_F\ x\ in\ at\ x0. b \in existence\text{-}ivl\ t0\ x$

<proof>

lemma *uniform-limit-flow:*

assumes $a: a \in existence\text{-}ivl\ t0\ x0$

assumes $b: b \in existence\text{-}ivl\ t0\ x0$

assumes $le: a \leq b$

shows $uniform\text{-}limit\ \{a .. b\}\ (flow\ t0)\ (flow\ t0\ x0)\ (at\ x0)$

<proof>

lemma *eventually-at-fst:*

assumes $eventually\ P\ (at\ (fst\ x))$

assumes $P\ (fst\ x)$

shows $eventually\ (\lambda h. P\ (fst\ h))\ (at\ x)$

<proof>

lemma *eventually-at-snd:*

assumes $eventually\ P\ (at\ (snd\ x))$

assumes $P\ (snd\ x)$

shows $eventually\ (\lambda h. P\ (snd\ h))\ (at\ x)$

<proof>

lemma

shows *open-state-space*: *open* (*Sigma X (existence-ivl t0)*)

and *flow-continuous-on-state-space*:

continuous-on (Sigma X (existence-ivl t0)) ($\lambda(x, t). \text{flow } t0 \ x \ t$)

<proof>

lemmas *flow-continuous-on-compose*[*continuous-intros*] =

continuous-on-compose-Pair[*OF flow-continuous-on-state-space*]

lemma *flow-isCont-state-space*: $t \in \text{existence-ivl } t0 \ x0 \implies \text{isCont } (\lambda(x, t). \text{flow } t0 \ x \ t) \ (x0, t)$

<proof>

lemma

flow-absolutely-integrable-on[*integrable-on-simps*]:

assumes $s \in \text{existence-ivl } t0 \ x0$

shows $(\lambda x. \text{norm } (\text{flow } t0 \ x0 \ x)) \text{ integrable-on closed-segment } t0 \ s$

<proof>

lemma *existence-ivl-eq-domain*:

assumes *iv-defined*: $t0 \in T \ x0 \in X$

assumes *bnd*: $\bigwedge tm \ tM \ t \ x. tm \in T \implies tM \in T \implies \exists M. \exists L. \forall t \in \{tm \ .. \ tM\}.$

$\forall x \in X. \text{norm } (f \ t \ x) \leq M + L * \text{norm } x$

assumes *is-interval* $T \ X = UNIV$

shows $\text{existence-ivl } t0 \ x0 = T$

<proof>

lemma *flow-unique*:

assumes $t \in \text{existence-ivl } t0 \ x0$

assumes $\text{phi } t0 = x0$

assumes $\bigwedge t. t \in \text{existence-ivl } t0 \ x0 \implies (\text{phi has-vector-derivative } f \ t \ (\text{phi } t))$

(*at t*)

assumes $\bigwedge t. t \in \text{existence-ivl } t0 \ x0 \implies \text{phi } t \in X$

shows $\text{flow } t0 \ x0 \ t = \text{phi } t$

<proof>

lemma *flow-unique-on*:

assumes $t \in \text{existence-ivl } t0 \ x0$

assumes $\text{phi } t0 = x0$

assumes (*phi has-vderiv-on* ($\lambda t. f \ t \ (\text{phi } t)$)) (*existence-ivl t0 x0*)

assumes $\bigwedge t. t \in \text{existence-ivl } t0 \ x0 \implies \text{phi } t \in X$

shows $\text{flow } t0 \ x0 \ t = \text{phi } t$

<proof>

end — *local-lipschitz T X f*

locale *two-ll-on-open* =

F: ll-on-open T1 F X + G: ll-on-open T2 G X
for F T1 G T2 X J x0 +
fixes e::real and K
assumes t0-in-J: 0 ∈ J
assumes J-subset: J ⊆ F.existence-ivl 0 x0
assumes J-ivl: is-interval J
assumes F-lipschitz: $\bigwedge t. t \in J \implies K\text{-lipschitz-on } X (F t)$
assumes K-pos: 0 < K
assumes F-G-norm-ineq: $\bigwedge t x. t \in J \implies x \in X \implies \text{norm } (F t x - G t x) < e$
begin

context begin

lemma F-iv-defined: 0 ∈ T1 x0 ∈ X
 ⟨proof⟩

lemma e-pos: 0 < e
 ⟨proof⟩ **definition flow0 t = F.flow 0 x0 t**
qualified definition Y t = G.flow 0 x0 t

lemma norm-X-Y-bound:
shows $\forall t \in J \cap G.\text{existence-ivl } 0 x0. \text{norm } (\text{flow0 } t - Y t) \leq e / K * (\exp(K * |t|) - 1)$
 ⟨proof⟩

end

end

locale auto-ll-on-open =
fixes f::'a::{banach, heine-borel} \Rightarrow 'a and X
assumes auto-local-lipschitz: local-lipschitz UNIV X (λ ::real. f)
assumes auto-open-domain[intro!, simp]: open X
begin

autonomous flow and existence interval

definition flow0 x0 t = ll-on-open.flow UNIV (λ -. f) X 0 x0 t

definition existence-ivl0 x0 = ll-on-open.existence-ivl UNIV (λ -. f) X 0 x0

sublocale ll-on-open-it UNIV λ -. f X 0
rewrites flow = ($\lambda t0 x0 t. \text{flow0 } x0 (t - t0)$)
and existence-ivl = ($\lambda t0 x0. (+) t0 \text{ 'existence-ivl0 } x0$)
and (+) 0 = ($\lambda x::\text{real}. x$)
and s - 0 = s
and ($\lambda x. x$) ' S = S
and s ∈ (+) t ' S \longleftrightarrow s - t ∈ (S::real set)
and P (s + t - s) = P (t::real) — TODO: why does just the equation not work?

and $P (t + s - s) = P t$ — TODO: why does just the equation not work?
 ⟨proof⟩

lemma *existence-ivl-zero*: $x0 \in X \implies 0 \in \text{existence-ivl0 } x0$ ⟨proof⟩

lemmas [*continuous-intros del*] = *continuous-on-f*

lemmas *continuous-on-f-comp*[*continuous-intros*] = *continuous-on-f*[*OF continuous-on-const - subset-UNIV*]

lemma

flow-in-compact-right-existence:

assumes $\bigwedge t. 0 \leq t \implies t \in \text{existence-ivl0 } x \implies \text{flow0 } x t \in K$

assumes *compact* $K K \subseteq X$

assumes $x \in X t \geq 0$

shows $t \in \text{existence-ivl0 } x$

⟨proof⟩

lemma

flow-in-compact-left-existence:

assumes $\bigwedge t. t \leq 0 \implies t \in \text{existence-ivl0 } x \implies \text{flow0 } x t \in K$

assumes *compact* $K K \subseteq X$

assumes $x \in X t \leq 0$

shows $t \in \text{existence-ivl0 } x$

⟨proof⟩

end

locale *compact-continuously-diff* =

derivative-on-prod $T X f \lambda(t, x). f' x o_L \text{snd-blinfun}$

for $T X$ **and** $f::\text{real} \Rightarrow 'a::\{\text{banach,perfect-space,heine-borel}\} \Rightarrow 'a$

and $f'::'a \Rightarrow ('a, 'a) \text{blinfun} +$

assumes *compact-domain*: *compact* X

assumes *convex*: *convex* X

assumes *nonempty-domains*: $T \neq \{\}$ $X \neq \{\}$

assumes *continuous-derivative*: *continuous-on* $X f'$

begin

lemma *ex-onorm-bound*:

$\exists B. \forall x \in X. \text{norm } (f' x) \leq B$

⟨proof⟩

definition *onorm-bound* = (*SOME* $B. \forall x \in X. \text{norm } (f' x) \leq B$)

lemma *onorm-bound*: **assumes** $x \in X$ **shows** $\text{norm } (f' x) \leq \text{onorm-bound}$

⟨proof⟩

sublocale *closed-domain* X

⟨proof⟩

```

sublocale global-lipschitz  $T X f$  onorm-bound
  ⟨proof⟩

end — compact  $X$ 

locale unique-on-compact-continuously-diff = self-mapping +
  compact-interval  $T$  +
  compact-continuously-diff  $T X f$ 
begin

sublocale unique-on-closed  $t0 T x0 f X$  onorm-bound
  ⟨proof⟩

end

locale c1-on-open =
  fixes  $f::'a::\{banach, perfect-space, heine-borel\} \Rightarrow 'a$  and  $f' X$ 
  assumes open-dom[simp]: open  $X$ 
  assumes derivative-rhs:
     $\bigwedge x. x \in X \implies (f \text{ has-derivative } blinfun-apply (f' x)) (at x)$ 
  assumes continuous-derivative: continuous-on  $X f'$ 
begin

lemmas continuous-derivative-comp[continuous-intros] =
  continuous-on-compose2[OF continuous-derivative]

lemma derivative-tendsto[tendsto-intros]:
  assumes [tendsto-intros]:  $(g \longrightarrow l) F$ 
  and  $l \in X$ 
  shows  $((\lambda x. f' (g x)) \longrightarrow f' l) F$ 
  ⟨proof⟩

lemma c1-on-open-rev[intro, simp]: c1-on-open  $(-f) (-f')$   $X$ 
  ⟨proof⟩

lemma derivative-rhs-compose[derivative-intros]:
   $((g \text{ has-derivative } g') (at x \text{ within } s)) \implies g x \in X \implies$ 
   $((\lambda x. f (g x)) \text{ has-derivative } (\lambda xa. blinfun-apply (f' (g x)) (g' xa)))$ 
   $(at x \text{ within } s)$ 
  ⟨proof⟩

sublocale auto-ll-on-open
  ⟨proof⟩

end —  $?x \in X \implies (f \text{ has-derivative } blinfun-apply (f' ?x)) (at ?x)$ 

locale c1-on-open-euclidean = c1-on-open  $f f' X$ 
  for  $f::'a::euclidean-space \Rightarrow -$  and  $f' X$ 

```

begin

lemma *c1-on-open-euclidean-anchor*: True \langle proof \rangle

definition *vareq* $x0\ t = f'$ (*flow0* $x0\ t$)

interpretation *var*: *ll-on-open-existence-ivl0* $x0\ vareq\ x0\ UNIV$
 \langle proof \rangle

context begin

lemma *continuous-on-A*[*continuous-intros*]:

assumes *continuous-on* $S\ a$

assumes *continuous-on* $S\ b$

assumes $\bigwedge s. s \in S \implies a\ s \in X$

assumes $\bigwedge s. s \in S \implies b\ s \in \textit{existence-ivl0}\ (a\ s)$

shows *continuous-on* $S\ (\lambda s. \textit{vareq}\ (a\ s)\ (b\ s))$

\langle proof \rangle

lemmas [*intro*] = *mem-existence-ivl-iv-defined*

context

fixes $x0::'a$

begin

lemma *flow0-defined*: $xa \in \textit{existence-ivl0}\ x0 \implies \textit{flow0}\ x0\ xa \in X$
 \langle proof \rangle

lemma *continuous-on-flow0*: *continuous-on* (*existence-ivl0* $x0$) (*flow0* $x0$)
 \langle proof \rangle

lemmas *continuous-on-flow0-comp*[*continuous-intros*] = *continuous-on-compose2*[*OF*
continuous-on-flow0]

lemma *varexivl-eq-exivl*:

assumes $t \in \textit{existence-ivl0}\ x0$

shows $\textit{var.existence-ivl}\ x0\ t\ a = \textit{existence-ivl0}\ x0$

\langle proof \rangle

definition *vector-Dflow* $u0\ t \equiv \textit{var.flow}\ x0\ 0\ u0\ t$

qualified abbreviation $Y\ z\ t \equiv \textit{flow0}\ (x0 + z)\ t$

Linearity of the solution to the variational equation. TODO: generalize this and some other things for arbitrary linear ODEs

lemma *vector-Dflow-linear*:

assumes $t \in \textit{existence-ivl0}\ x0$

shows $\textit{vector-Dflow}\ (\alpha *_{\mathbb{R}} a + \beta *_{\mathbb{R}} b)\ t = \alpha *_{\mathbb{R}} \textit{vector-Dflow}\ a\ t + \beta *_{\mathbb{R}} \textit{vector-Dflow}\ b\ t$

\langle proof \rangle

lemma *linear-vector-Dflow*:

assumes $t \in \text{existence-ivl0 } x0$

shows *linear* $(\lambda z. \text{vector-Dflow } z \ t)$

$\langle \text{proof} \rangle$

lemma *bounded-linear-vector-Dflow*:

assumes $t \in \text{existence-ivl0 } x0$

shows *bounded-linear* $(\lambda z. \text{vector-Dflow } z \ t)$

$\langle \text{proof} \rangle$

lemma *vector-Dflow-continuous-on-time*: $x0 \in X \implies \text{continuous-on } (\text{existence-ivl0 } x0)$ $(\lambda t. \text{vector-Dflow } z \ t)$

$\langle \text{proof} \rangle$

proposition *proposition-17-6-weak*:

— from "Differential Equations, Dynamical Systems, and an Introduction to Chaos", Hirsch/Smale/Devaney

assumes $t \in \text{existence-ivl0 } x0$

shows $(\lambda y. (Y (y - x0) \ t - \text{flow0 } x0 \ t - \text{vector-Dflow } (y - x0) \ t) /_R \text{norm } (y - x0)) - x0 \rightarrow 0$

$\langle \text{proof} \rangle$

lemma *local-lipschitz-A*:

$OT \subseteq \text{existence-ivl0 } x0 \implies \text{local-lipschitz } OT \ (\text{OS}::('a \Rightarrow_L 'a) \ \text{set}) \ (\lambda t. (o_L) \ (\text{vareq } x0 \ t))$

$\langle \text{proof} \rangle$

lemma *total-derivative-ll-on-open*:

ll-on-open $(\text{existence-ivl0 } x0)$ $(\lambda t. \text{blinfun-compose } (\text{vareq } x0 \ t)) \ (\text{UNIV}::('a \Rightarrow_L 'a) \ \text{set})$

$\langle \text{proof} \rangle$

end

end

sublocale *mvar*: *ll-on-open existence-ivl0 x0* $\lambda t. \text{blinfun-compose } (\text{vareq } x0 \ t) \ \text{UNIV}::('a \Rightarrow_L 'a) \ \text{set}$ **for** $x0$

$\langle \text{proof} \rangle$

lemma *mvar-existence-ivl-eq-existence-ivl[simp]*:— TODO: unify with $?t \in \text{existence-ivl0 } ?x0.0 \implies \text{var.existence-ivl } ?x0.0 \ ?t \ ?a = \text{existence-ivl0 } ?x0.0$

assumes $t \in \text{existence-ivl0 } x0$

shows $\text{mvar.existence-ivl } x0 \ t = (\lambda-. \text{existence-ivl0 } x0)$

$\langle \text{proof} \rangle$

lemma

assumes $t \in \text{existence-ivl0 } x0$

shows *continuous-on* ($UNIV \times \text{existence-ivl0 } x0$) $(\lambda(x, ta). \text{mvar.flow } x0 \ t \ x \ ta)$
 $\langle \text{proof} \rangle$

definition $D\text{flow } x0 = \text{mvar.flow } x0 \ 0 \ \text{id-blinfun}$

lemma *var-eq-mvar*:

assumes $t0 \in \text{existence-ivl0 } x0$

assumes $t \in \text{existence-ivl0 } x0$

shows $\text{var.flow } x0 \ t0 \ i \ t = \text{mvar.flow } x0 \ t0 \ \text{id-blinfun } t \ i$

$\langle \text{proof} \rangle$

lemma *Dflow-zero[simp]*: $x \in X \implies D\text{flow } x \ 0 = 1_L$

$\langle \text{proof} \rangle$

5.3 Differentiability of the flow0

$U \ t$, i.e. the solution of the variational equation, is the space derivative at the initial value $x0$.

lemma *flow-dx-derivative*:

assumes $t \in \text{existence-ivl0 } x0$

shows $((\lambda x0. \text{flow0 } x0 \ t) \text{ has-derivative } (\lambda z. \text{vector-Dflow } x0 \ z \ t)) \text{ (at } x0)$

$\langle \text{proof} \rangle$

lemma *flow-dx-derivative-blinfun*:

assumes $t \in \text{existence-ivl0 } x0$

shows $((\lambda x. \text{flow0 } x \ t) \text{ has-derivative } \text{Blinfun } (\lambda z. \text{vector-Dflow } x0 \ z \ t)) \text{ (at } x0)$

$\langle \text{proof} \rangle$

definition $\text{floweriv } x0 \ t = \text{comp12 } (D\text{flow } x0 \ t) \ (\text{blinfun-scaleR-left } (f \ (\text{flow0 } x0 \ t)))$

lemma *floweriv-eq*: $\text{floweriv } x0 \ t \ (\xi_1, \xi_2) = (D\text{flow } x0 \ t) \ \xi_1 + \xi_2 *_R f \ (\text{flow0 } x0 \ t)$

$\langle \text{proof} \rangle$

lemma *W-continuous-on*: *continuous-on* ($\text{Sigma } X \ \text{existence-ivl0}$) $(\lambda(x0, t). D\text{flow } x0 \ t)$

— TODO: somewhere here is hidden continuity wrt rhs of ODE, extract it!

$\langle \text{proof} \rangle$

lemma *W-continuous-on-comp[continuous-intros]*:

assumes $h: \text{continuous-on } S \ h$ **and** $g: \text{continuous-on } S \ g$

shows $(\bigwedge s. s \in S \implies h \ s \in X) \implies (\bigwedge s. s \in S \implies g \ s \in \text{existence-ivl0 } (h \ s))$

\implies

$\text{continuous-on } S \ (\lambda s. D\text{flow } (h \ s) \ (g \ s))$

$\langle \text{proof} \rangle$

lemma *f-flow-continuous-on*: *continuous-on* ($\text{Sigma } X \ \text{existence-ivl0}$) $(\lambda(x0, t). f \ (\text{flow0 } x0 \ t))$

<proof>

lemma

flow-has-space-derivative:

assumes $t \in \text{existence-ivl0 } x0$

shows $((\lambda x0. \text{flow0 } x0 t) \text{ has-derivative } D\text{flow } x0 t) \text{ (at } x0)$

<proof>

lemma

flow-has-flowderiv:

assumes $t \in \text{existence-ivl0 } x0$

shows $((\lambda(x0, t). \text{flow0 } x0 t) \text{ has-derivative } \text{flowerderiv } x0 t) \text{ (at } (x0, t) \text{ within } S)$

<proof>

lemma *flow0-comp-has-derivative:*

assumes $h: h s \in \text{existence-ivl0 } (g s)$

assumes $[\text{derivative-intros}]: (g \text{ has-derivative } g') \text{ (at } s \text{ within } S)$

assumes $[\text{derivative-intros}]: (h \text{ has-derivative } h') \text{ (at } s \text{ within } S)$

shows $((\lambda x. \text{flow0 } (g x) (h x)) \text{ has-derivative } (\lambda x. \text{blinfun-apply } (\text{flowerderiv } (g s) (h s)) (g' x, h' x)))$

(at } s \text{ within } S)

<proof>

lemma *flowerderiv-continuous-on: continuous-on* $(\text{Sigma } X \text{ existence-ivl0}) (\lambda(x0, t). \text{flowerderiv } x0 t)$

<proof>

lemma *flowerderiv-continuous-on-comp* $[\text{continuous-intros}]$:

assumes *continuous-on* $S x$

assumes *continuous-on* $S t$

assumes $\bigwedge s. s \in S \implies x s \in X \bigwedge s. s \in S \implies t s \in \text{existence-ivl0 } (x s)$

shows *continuous-on* $S (\lambda xa. \text{flowerderiv } (x xa) (t xa))$

<proof>

lemmas $[\text{intro}] = \text{flow-in-domain}$

lemma *vareq-trans*: $t0 \in \text{existence-ivl0 } x0 \implies t \in \text{existence-ivl0 } (\text{flow0 } x0 t0) \implies$

$\text{vareq } (\text{flow0 } x0 t0) t = \text{vareq } x0 (t0 + t)$

<proof>

lemma *diff-existence-ivl-trans*:

$t0 \in \text{existence-ivl0 } x0 \implies t \in \text{existence-ivl0 } x0 \implies t - t0 \in \text{existence-ivl0 } (\text{flow0 } x0 t0)$ **for** t

<proof>

lemma *has-vderiv-on-blinfun-compose-right* $[\text{derivative-intros}]$:

assumes $(g \text{ has-vderiv-on } g') T$

assumes $\bigwedge x. x \in T \implies gd' x = g' x o_L d$

shows $((\lambda x. g x o_L d) \text{ has-vderiv-on } gd') T$

<proof>

lemma *has-vderiv-on-blinfun-compose-left*[*derivative-intros*]:

assumes (*g has-vderiv-on g'*) *T*
assumes $\bigwedge x. x \in T \implies g d' x = d \circ_L g' x$
shows $((\lambda x. d \circ_L g x) \text{ has-vderiv-on } g d') T$
<proof>

lemma *mvar-flow-shift*:

assumes *t0* \in *existence-ivl0* *x0* *t1* \in *existence-ivl0* *x0*
shows *mvar.flow* *x0* *t0* *d* *t1* = *Dflow* (*flow0* *x0* *t0*) (*t1* - *t0*) \circ_L *d*
<proof>

lemma *Dflow-trans*:

assumes *h* \in *existence-ivl0* *x0*
assumes *i* \in *existence-ivl0* (*flow0* *x0* *h*)
shows *Dflow* *x0* (*h* + *i*) = *Dflow* (*flow0* *x0* *h*) *i* \circ_L (*Dflow* *x0* *h*)
<proof>

lemma *Dflow-trans-apply*:

assumes *h* \in *existence-ivl0* *x0*
assumes *i* \in *existence-ivl0* (*flow0* *x0* *h*)
shows *Dflow* *x0* (*h* + *i*) *d0* = *Dflow* (*flow0* *x0* *h*) *i* (*Dflow* *x0* *h* *d0*)
<proof>

end — *True*

end

6 Upper and Lower Solutions

theory *Upper-Lower-Solution*

imports *Flow*

begin

Following Walter [1] in section 9

lemma *IVT-min*:

fixes *f* :: *real* \Rightarrow '*b* :: {*linorder-topology, real-normed-vector, ordered-real-vector*}
— generalize?
assumes *y*: *f* *a* \leq *y* *y* \leq *f* *b* *a* \leq *b*
assumes *: *continuous-on* {*a* .. *b*} *f*
notes [*continuous-intros*] = * [*THEN* *continuous-on-subset*]
obtains *x* **where** *a* \leq *x* *x* \leq *b* *f* *x* = *y* $\bigwedge x'. a \leq x' \implies x' < x \implies f x' < y$
<proof>

lemma *filtermap-at-left-shift*: *filtermap* ($\lambda x. x - d$) (*at-left* *a*) = *at-left* (*a* - *d*::*real*)

<proof>

context

fixes $v v' w w' :: \text{real} \Rightarrow \text{real}$ **and** $t0 t1 e :: \text{real}$
assumes v' : (v has-vderiv-on v') $\{t0 <.. t1\}$
and w' : (w has-vderiv-on w') $\{t0 <.. t1\}$
assumes $pos-ivl$: $t0 < t1$
assumes $e-pos$: $e > 0$ **and** $e-in$: $t0 + e \leq t1$
assumes $less$: $\bigwedge t. t0 < t \implies t < t0 + e \implies v t < w t$
begin

lemma *first-intersection-crossing-derivatives*:

assumes na : $t0 < tg \leq t1$ $v tg \geq w tg$
notes [*continuous-intros*] =
 $vderiv-on-continuous-on[OF v', THEN continuous-on-subset]$
 $vderiv-on-continuous-on[OF w', THEN continuous-on-subset]$
obtains $x0$ **where**
 $t0 < x0 \leq tg$
 $v' x0 \geq w' x0$
 $v x0 = w x0$
 $\bigwedge t. t0 < t \implies t < x0 \implies v t < w t$
<proof>

lemma *defect-less*:

assumes b : $\bigwedge t. t0 < t \implies t \leq t1 \implies v' t - f t (v t) < w' t - f t (w t)$
notes [*continuous-intros*] =
 $vderiv-on-continuous-on[OF v', THEN continuous-on-subset]$
 $vderiv-on-continuous-on[OF w', THEN continuous-on-subset]$
shows $\forall t \in \{t0 <.. t1\}. v t < w t$
<proof>

end

lemma *has-derivatives-less-lemma*:

fixes $v v' :: \text{real} \Rightarrow \text{real}$
assumes v' : (v has-vderiv-on v') T
assumes y' : (y has-vderiv-on y') T
assumes lu : $\bigwedge t. t \in T \implies t > t0 \implies v' t - f t (v t) < y' t - f t (y t)$
assumes $lower$: $v t0 \leq y t0$
assumes $eq-imp$: $v t0 = y t0 \implies v' t0 < y' t0$
assumes t : $t0 < t$ $t0 \in T$ $t \in T$ *is-interval* T
shows $v t < y t$
<proof>

lemma *strict-lower-solution*:

fixes $v v' :: \text{real} \Rightarrow \text{real}$
assumes sol : (y solves-ode f) $T X$
assumes v' : (v has-vderiv-on v') T
assumes $lower$: $\bigwedge t. t \in T \implies t > t0 \implies v' t < f t (v t)$
assumes iv : $v t0 \leq y t0$ $v t0 = y t0 \implies v' t0 < f t0 (y t0)$

assumes $t: t0 < t \ t0 \in T \ t \in T \text{ is-interval } T$
shows $v \ t < y \ t$
 <proof>

lemma *strict-upper-solution*:

fixes $w \ w': \text{real} \Rightarrow \text{real}$
assumes $\text{sol}: (y \text{ solves-ode } f) \ T \ X$
assumes $w': (w \text{ has-vderiv-on } w') \ T$
and $\text{upper}: \bigwedge t. t \in T \Longrightarrow t > t0 \Longrightarrow f \ t \ (w \ t) < w' \ t$
and $\text{iv}: y \ t0 \leq w \ t0 \ y \ t0 = w \ t0 \Longrightarrow f \ t0 \ (y \ t0) < w' \ t0$
assumes $t: t0 < t \ t0 \in T \ t \in T \text{ is-interval } T$
shows $y \ t < w \ t$
 <proof>

lemma *uniform-limit-at-within-subset*:

assumes $\text{uniform-limit } S \ x \ l \ (\text{at } t \ \text{within } T)$
assumes $U \subseteq T$
shows $\text{uniform-limit } S \ x \ l \ (\text{at } t \ \text{within } U)$
 <proof>

lemma *uniform-limit-le*:

fixes $f::'c \Rightarrow 'a \Rightarrow 'b::\{\text{metric-space, linorder-topology}\}$
assumes $I: I \neq \text{bot}$
assumes $u: \text{uniform-limit } X \ f \ g \ I$
assumes $u': \text{uniform-limit } X \ f' \ g' \ I$
assumes $\forall_F \ i \ \text{in } I. \forall x \in X. f \ i \ x \leq f' \ i \ x$
assumes $x \in X$
shows $g \ x \leq g' \ x$
 <proof>

lemma *uniform-limit-le-const*:

fixes $f::'c \Rightarrow 'a \Rightarrow 'b::\{\text{metric-space, linorder-topology}\}$
assumes $I: I \neq \text{bot}$
assumes $u: \text{uniform-limit } X \ f \ g \ I$
assumes $\forall_F \ i \ \text{in } I. \forall x \in X. f \ i \ x \leq h \ x$
assumes $x \in X$
shows $g \ x \leq h \ x$
 <proof>

lemma *uniform-limit-ge-const*:

fixes $f::'c \Rightarrow 'a \Rightarrow 'b::\{\text{metric-space, linorder-topology}\}$
assumes $I: I \neq \text{bot}$
assumes $u: \text{uniform-limit } X \ f \ g \ I$
assumes $\forall_F \ i \ \text{in } I. \forall x \in X. h \ x \leq f \ i \ x$
assumes $x \in X$
shows $h \ x \leq g \ x$
 <proof>

locale *ll-on-open-real* = *ll-on-open* $T \ f \ X$ **for** $T \ f$ **and** $X::\text{real set}$

begin

lemma *lower-solution*:

fixes $v\ v' :: \text{real} \Rightarrow \text{real}$

assumes $\text{sol}: (y \text{ solves-ode } f) S\ X$

assumes $v': (v \text{ has-vderiv-on } v') S$

assumes $\text{lower}: \bigwedge t. t \in S \implies t > t0 \implies v' t < f t (v t)$

assumes $\text{iv}: v\ t0 \leq y\ t0$

assumes $t: t0 \leq t\ t0 \in S\ t \in S \text{ is-interval } S\ S \subseteq T$

shows $v\ t \leq y\ t$

<proof>

lemma *upper-solution*:

fixes $v\ v' :: \text{real} \Rightarrow \text{real}$

assumes $\text{sol}: (y \text{ solves-ode } f) S\ X$

assumes $v': (v \text{ has-vderiv-on } v') S$

assumes $\text{upper}: \bigwedge t. t \in S \implies t > t0 \implies f t (v t) < v' t$

assumes $\text{iv}: y\ t0 \leq v\ t0$

assumes $t: t0 \leq t\ t0 \in S\ t \in S \text{ is-interval } S\ S \subseteq T$

shows $y\ t \leq v\ t$

<proof>

end

end

theory *Poincare-Map*

imports

Flow

begin

abbreviation $\text{plane } n\ c \equiv \{x. x \cdot n = c\}$

lemma

eventually-tendsto-compose-within:

assumes $\text{eventually } P \text{ (at } l \text{ within } S)$

assumes $P\ l$

assumes $(f \longrightarrow l) \text{ (at } x \text{ within } T)$

assumes $\text{eventually } (\lambda x. f\ x \in S) \text{ (at } x \text{ within } T)$

shows $\text{eventually } (\lambda x. P (f\ x)) \text{ (at } x \text{ within } T)$

<proof>

lemma

eventually-eventually-withinI:—aha...

assumes $\forall_F x \text{ in at } x \text{ within } A. P\ x\ P\ x$

shows $\forall_F a \text{ in at } x \text{ within } S. \forall_F x \text{ in at } a \text{ within } A. P\ x$

<proof>

lemma *eventually-not-in-closed*:

assumes $\text{closed } P$

assumes $f t \notin P \ t \in T$
assumes *continuous-on* $T f$
shows $\forall_F t$ in *at* t within T . $f t \notin P$
<proof>

context *ll-on-open-it* **begin**

lemma

existence-ivl-trans':
assumes $t + s \in \text{existence-ivl } t0 \ x0$
 $t \in \text{existence-ivl } t0 \ x0$
shows $t + s \in \text{existence-ivl } t \ (\text{flow } t0 \ x0 \ t)$
<proof>

end

context *auto-ll-on-open*— **TODO**: generalize to continuous systems
begin

definition *returns-to* $:: 'a \text{ set} \Rightarrow 'a \Rightarrow \text{bool}$

where *returns-to* $P \ x \longleftrightarrow (\forall_F t$ in *at-right* 0 . $\text{flow0 } x \ t \notin P) \wedge (\exists t > 0$. $t \in \text{existence-ivl0 } x \wedge \text{flow0 } x \ t \in P)$

definition *return-time* $:: 'a \text{ set} \Rightarrow 'a \Rightarrow \text{real}$

where *return-time* $P \ x =$
(if *returns-to* $P \ x$ *then* *(SOME* t .
 $t > 0 \wedge$
 $t \in \text{existence-ivl0 } x \wedge$
 $\text{flow0 } x \ t \in P \wedge$
 $(\forall s \in \{0 <..<t\}$. $\text{flow0 } x \ s \notin P))$ *else* $0)$

lemma *returns-toI*:

assumes $t: t > 0 \ t \in \text{existence-ivl0 } x \ \text{flow0 } x \ t \in P$
assumes *ev*: $\forall_F t$ in *at-right* 0 . $\text{flow0 } x \ t \notin P$
assumes *closed* P
shows *returns-to* $P \ x$
<proof>

lemma *returns-to-outsideI*:

assumes $t: t \geq 0 \ t \in \text{existence-ivl0 } x \ \text{flow0 } x \ t \in P$
assumes *ev*: $x \notin P$
assumes *closed* P
shows *returns-to* $P \ x$
<proof>

lemma *returns-toE*:

assumes *returns-to* $P \ x$
obtains $t0 \ t1$ **where**
 $0 < t0$

$t0 \leq t1$
 $t1 \in \text{existence-ivl0 } x$
 $\text{flow0 } x \ t1 \in P$
 $\bigwedge t. 0 < t \implies t < t0 \implies \text{flow0 } x \ t \notin P$
 <proof>

lemma *return-time-some*:
assumes *returns-to P x*
shows *return-time P x =*
 (*SOME t. t > 0 \wedge t \in existence-ivl0 x \wedge flow0 x t \in P \wedge ($\forall s \in \{0 <..<t\}$.
 flow0 x s \notin P))*
 <proof>

lemma *return-time-ex1*:
assumes *returns-to P x*
assumes *closed P*
shows $\exists! t. t > 0 \wedge t \in \text{existence-ivl0 } x \wedge \text{flow0 } x \ t \in P \wedge (\forall s \in \{0 <..<t\}.$
 $\text{flow0 } x \ s \notin P)$
 <proof>

lemma
return-time-pos-returns-to:
return-time P x > 0 \implies returns-to P x
 <proof>

lemma
assumes *ret: returns-to P x*
assumes *closed P*
shows *return-time-pos: return-time P x > 0*
 <proof>

lemma *returns-to-return-time-pos*:
assumes *closed P*
shows *returns-to P x \longleftrightarrow return-time P x > 0*
 <proof>

lemma *return-time*:
assumes *ret: returns-to P x*
assumes *closed P*
shows *return-time P x > 0*
and *return-time-exivl: return-time P x \in existence-ivl0 x*
and *return-time-returns: flow0 x (return-time P x) \in P*
and *return-time-least: $\bigwedge s. 0 < s \implies s < \text{return-time P x} \implies \text{flow0 } x \ s \notin P$*
 <proof>

lemma *returns-to-earlierI*:
assumes *ret: returns-to P (flow0 x t) closed P*
assumes *t \geq 0 t \in existence-ivl0 x*
assumes *ev: $\forall_F t$ in at-right 0. flow0 x t \notin P*

shows *returns-to* P x
<proof>

lemma *return-time-gt*:

assumes *ret*: *returns-to* P x *closed* P
assumes *flow-not*: $\bigwedge s. 0 < s \implies s \leq t \implies \text{flow0 } x \ s \notin P$
shows $t < \text{return-time } P \ x$
<proof>

lemma *return-time-le*:

assumes *ret*: *returns-to* P x *closed* P
assumes *flow-not*: $\text{flow0 } x \ t \in P \ t > 0$
shows $\text{return-time } P \ x \leq t$
<proof>

lemma *returns-to-laterI*:

assumes *ret*: *returns-to* P x *closed* P
assumes *t*: $t > 0 \ t \in \text{existence-ivl0 } x$
assumes *flow-not*: $\bigwedge s. 0 < s \implies s \leq t \implies \text{flow0 } x \ s \notin P$
shows *returns-to* P $(\text{flow0 } x \ t)$
<proof>

lemma *never-returns*:

assumes $\neg \text{returns-to } P \ x$
assumes *closed* $P \ t \geq 0 \ t \in \text{existence-ivl0 } x$
assumes *ev*: $\forall_F \ t \text{ in at-right } 0. \ \text{flow0 } x \ t \notin P$
shows $\neg \text{returns-to } P \ (\text{flow0 } x \ t)$
<proof>

lemma *return-time-eqI*:

assumes *closed* P
and *t-pos*: $t > 0$
and *ex*: $t \in \text{existence-ivl0 } x$
and *ret*: $\text{flow0 } x \ t \in P$
and *least*: $\bigwedge s. 0 < s \implies s < t \implies \text{flow0 } x \ s \notin P$
shows $\text{return-time } P \ x = t$
<proof>

lemma *return-time-step*:

assumes *returns-to* P $(\text{flow0 } x \ t)$
assumes *closed* P
assumes *flow-not*: $\bigwedge s. 0 < s \implies s \leq t \implies \text{flow0 } x \ s \notin P$
assumes *t*: $t > 0 \ t \in \text{existence-ivl0 } x$
shows $\text{return-time } P \ (\text{flow0 } x \ t) = \text{return-time } P \ x - t$
<proof>

definition *poincare-map* $P \ x = \text{flow0 } x \ (\text{return-time } P \ x)$

lemma *poincare-map-step-flow*:

assumes *ret*: returns-to P x closed P
assumes *flow-not*: $\bigwedge s. 0 < s \implies s \leq t \implies \text{flow0 } x \ s \notin P$
assumes $t: t > 0 \ t \in \text{existence-ivl0 } x$
shows $\text{poincare-map } P (\text{flow0 } x \ t) = \text{poincare-map } P \ x$
 $\langle \text{proof} \rangle$

lemma *poincare-map-returns*:
assumes returns-to P x closed P
shows $\text{poincare-map } P \ x \in P$
 $\langle \text{proof} \rangle$

lemma *poincare-map-onto*:
assumes closed P
assumes $0 < t \ t \in \text{existence-ivl0 } x \ \forall_F \ t \ \text{in at-right } 0. \ \text{flow0 } x \ t \notin P$
assumes $\text{flow0 } x \ t \in P$
shows $\text{poincare-map } P \ x \in \text{flow0 } x \ ' \{0 <.. t\} \cap P$
 $\langle \text{proof} \rangle$

end

lemma *isCont-blinfunD*:
fixes $f': 'a::\text{metric-space} \Rightarrow 'b::\text{real-normed-vector} \Rightarrow_L 'c::\text{real-normed-vector}$
assumes $\text{isCont } f' \ a \ 0 < e$
shows $\exists d > 0. \ \forall x. \ \text{dist } a \ x < d \longrightarrow \text{onorm } (\lambda v. \ \text{blinfun-apply } (f' \ x) \ v) - \text{blin-}$
 $\text{fun-apply } (f' \ a) \ v < e$
 $\langle \text{proof} \rangle$

proposition *has-derivative-locally-injective-blinfun*:
fixes $f :: 'n::\text{euclidean-space} \Rightarrow 'm::\text{euclidean-space}$
and $f': 'n \Rightarrow 'n \Rightarrow_L 'm$
and $g': 'm \Rightarrow_L 'n$
assumes $a \in s$
and open s
and $g': g' \ o_L (f' \ a) = 1_L$
and $f': \bigwedge x. \ x \in s \implies (f \ \text{has-derivative } f' \ x) \ (\text{at } x)$
and $c: \text{isCont } f' \ a$
obtains r **where** $r > 0 \ \text{ball } a \ r \subseteq s \ \text{inj-on } f \ (\text{ball } a \ r)$
 $\langle \text{proof} \rangle$

lift-definition $\text{embed1-blinfun}:: 'a::\text{real-normed-vector} \Rightarrow_L ('a * 'b::\text{real-normed-vector})$
is $\lambda x. (x, 0)$
 $\langle \text{proof} \rangle$

lemma *blinfun-apply-embed1-blinfun[simp]*: $\text{blinfun-apply } \text{embed1-blinfun} \ x = (x, 0)$
 $\langle \text{proof} \rangle$

lift-definition $\text{embed2-blinfun}:: 'a::\text{real-normed-vector} \Rightarrow_L ('b::\text{real-normed-vector} * 'a)$
is $\lambda x. (0, x)$

<proof>

lemma *blinfun-apply-embed2-blinfun[simp]*: *blinfun-apply embed2-blinfun* $x = (0, x)$

<proof>

lemma *blinfun-inverseD*: $f \circ_L f' = 1_L \implies f (f' x) = x$

<proof>

lemmas *continuous-on-open-vimageI = continuous-on-open-vimage*[*THEN iffD1, rule-format*]

lemmas *continuous-on-closed-vimageI = continuous-on-closed-vimage*[*THEN iffD1, rule-format*]

lemma *ball-times-subset*: $\text{ball } a \ (c/2) \times \text{ball } b \ (c/2) \subseteq \text{ball } (a, b) \ c$

<proof>

lemma *linear-inverse-blinop-lemma*:

fixes $w::'a::\{\text{banach, perfect-space}\}$ *blinop*

assumes $\text{norm } w < 1$

shows

summable $(\lambda n. (-1)^{\widehat{n}} *_{\mathbb{R}} w^{\widehat{n}})$ (**is** ?*C*)

$(\sum n. (-1)^{\widehat{n}} *_{\mathbb{R}} w^{\widehat{n}}) * (1 + w) = 1$ (**is** ?*I1*)

$(1 + w) * (\sum n. (-1)^{\widehat{n}} *_{\mathbb{R}} w^{\widehat{n}}) = 1$ (**is** ?*I2*)

$\text{norm } ((\sum n. (-1)^{\widehat{n}} *_{\mathbb{R}} w^{\widehat{n}}) - 1 + w) \leq (\text{norm } w)^2 / (1 - \text{norm } (w))$ (**is** ?*L*)

<proof>

lemma *linear-inverse-blinfun-lemma*:

fixes $w::'a \Rightarrow_L 'a::\{\text{banach, perfect-space}\}$

assumes $\text{norm } w < 1$

obtains *I* **where**

$I \circ_L (1_L + w) = 1_L (1_L + w) \circ_L I = 1_L$

$\text{norm } (I - 1_L + w) \leq (\text{norm } w)^2 / (1 - \text{norm } (w))$

<proof>

definition *invertibles-blinfun* = $\{w. \exists wi. w \circ_L wi = 1_L \wedge wi \circ_L w = 1_L\}$

lemma *blinfun-inverse-open*:— 8.3.2 in Dieudonne, TODO: add continuity and derivative

shows *open* (*invertibles-blinfun*::

$'a::\{\text{banach, perfect-space}\} \Rightarrow_L 'b::\text{banach}$) *set*)

<proof>

lemma *blinfun-compose-assoc[ac-simps]*: $a \circ_L b \circ_L c = a \circ_L (b \circ_L c)$

<proof>

TODO: move $\text{norm } (- ?x) = \text{norm } ?x$ to class!

lemma (**in** *real-normed-vector*) *norm-minus-cancel* [*simp*]: $\text{norm } (- x) = \text{norm } x$

<proof>

TODO: move $\text{norm } (?a - ?b) = \text{norm } (?b - ?a)$ to class!

lemma (in *real-normed-vector*) *norm-minus-commute*: $\text{norm } (a - b) = \text{norm } (b - a)$

<proof>

instance *euclidean-space* \subseteq *banach*

<proof>

lemma *blinfun-apply-Pair-split*:

$\text{blinfun-apply } g \ (a, b) = \text{blinfun-apply } g \ (a, 0) + \text{blinfun-apply } g \ (0, b)$

<proof>

lemma *blinfun-apply-Pair-add2*: $\text{blinfun-apply } f \ (0, a + b) = \text{blinfun-apply } f \ (0, a) + \text{blinfun-apply } f \ (0, b)$

<proof>

lemma *blinfun-apply-Pair-add1*: $\text{blinfun-apply } f \ (a + b, 0) = \text{blinfun-apply } f \ (a, 0) + \text{blinfun-apply } f \ (b, 0)$

<proof>

lemma *blinfun-apply-Pair-minus2*: $\text{blinfun-apply } f \ (0, a - b) = \text{blinfun-apply } f \ (0, a) - \text{blinfun-apply } f \ (0, b)$

<proof>

lemma *blinfun-apply-Pair-minus1*: $\text{blinfun-apply } f \ (a - b, 0) = \text{blinfun-apply } f \ (a, 0) - \text{blinfun-apply } f \ (b, 0)$

<proof>

lemma *implicit-function-theorem*:

fixes $f::'a::\text{euclidean-space} * 'b::\text{euclidean-space} \Rightarrow 'c::\text{euclidean-space}$ — **TODO**: generalize?!

assumes [*derivative-intros*]: $\bigwedge x. x \in S \implies (f \text{ has-derivative } \text{blinfun-apply } (f' \ x))$ (at x)

assumes $S: (x, y) \in S$ open S

assumes $\text{DIM}('c) \leq \text{DIM}('b)$

assumes $f'C: \text{isCont } f' \ (x, y)$

assumes $f \ (x, y) = 0$

assumes $T2: T \ o_L \ (f' \ (x, y) \ o_L \ \text{embed2-blinfun}) = 1_L$

assumes $T1: (f' \ (x, y) \ o_L \ \text{embed2-blinfun}) \ o_L \ T = 1_L$ — **TODO**: reduce?!

obtains $u \ e \ r$

where $f \ (x, u \ x) = 0$ $u \ x = y$

$\bigwedge s. s \in \text{cball } x \ e \implies f \ (s, u \ s) = 0$

continuous-on $(\text{cball } x \ e) \ u$

$(\lambda t. (t, u \ t)) \ ' \ \text{cball } x \ e \subseteq S$

$e > 0$

$(u \ \text{has-derivative} \ - \ T \ o_L \ f' \ (x, y) \ o_L \ \text{embed1-blinfun}) \ (\text{at } x)$

$r > 0$

$\bigwedge U \ v \ s. v \ x = y \implies (\bigwedge s. s \in U \implies f \ (s, v \ s) = 0) \implies U \subseteq \text{cball } x \ e \implies$

continuous-on $U \ v \implies s \in U \implies (s, v \ s) \in \text{ball } (x, y) \ r \implies u \ s = v \ s$

<proof>

lemma *implicit-function-theorem-unique:*

fixes $f::'a::\text{euclidean-space} * 'b::\text{euclidean-space} \Rightarrow 'c::\text{euclidean-space}$ — **TODO:** generalize?!

assumes $f'[derivative-intros]: \bigwedge x. x \in S \implies (f \text{ has-derivative } \text{blinfun-apply } (f' x)) \text{ (at } x)$

assumes $S: (x, y) \in S \text{ open } S$

assumes $D: DIM('c) \leq DIM('b)$

assumes $f'C: \text{continuous-on } S f'$

assumes $z: f(x, y) = 0$

assumes $T2: T \text{ o}_L (f'(x, y) \text{ o}_L \text{ embed2-blinfun}) = 1_L$

assumes $T1: (f'(x, y) \text{ o}_L \text{ embed2-blinfun}) \text{ o}_L T = 1_L$ — **TODO:** reduce?!

obtains $u e$

where $f(x, u x) = 0 \text{ u } x = y$

$\bigwedge s. s \in \text{cball } x e \implies f(s, u s) = 0$

$\text{continuous-on } (\text{cball } x e) u$

$(\lambda t. (t, u t)) ' \text{cball } x e \subseteq S$

$e > 0$

$(u \text{ has-derivative } (- T \text{ o}_L f'(x, y) \text{ o}_L \text{ embed1-blinfun})) \text{ (at } x)$

$\bigwedge s. s \in \text{cball } x e \implies f'(s, u s) \text{ o}_L \text{ embed2-blinfun} \in \text{invertibles-blinfun}$

$\bigwedge U v s. (\bigwedge s. s \in U \implies f(s, v s) = 0) \implies$

$u x = v x \implies$

$\text{continuous-on } U v \implies s \in U \implies x \in U \implies U \subseteq \text{cball } x e \implies \text{connected } U$

$\implies \text{open } U \implies u s = v s$

<proof>

lemma *uniform-limit-compose:*

assumes $ul: \text{uniform-limit } T f l F$

assumes $uc: \text{uniformly-continuous-on } S s$

assumes $ev: \forall_F x \text{ in } F. f x ' T \subseteq S$

assumes $subs: l ' T \subseteq S$

shows $\text{uniform-limit } T (\lambda i x. s (f i x)) (\lambda x. s (l x)) F$

<proof>

lemma

uniform-limit-in-open:

fixes $l::'a::\text{topological-space} \Rightarrow 'b::\text{heine-borel}$

assumes $ul: \text{uniform-limit } T f l \text{ (at } x)$

assumes $cont: \text{continuous-on } T l$

assumes $compact: \text{compact } T \text{ and } T\text{-ne: } T \neq \{\}$

assumes $B: \text{open } B$

assumes $mem: l ' T \subseteq B$

shows $\forall_F y \text{ in at } x. \forall t \in T. f y t \in B$

<proof>

lemma

order-uniform-limitD1:

fixes $l::'a::\text{topological-space} \Rightarrow \text{real}$ — **TODO:** generalize?!

assumes *ul*: *uniform-limit* T f l (at x)
assumes *cont*: *continuous-on* T l
assumes *compact*: *compact* T
assumes *less*: $\bigwedge t. t \in T \implies l\ t < b$
shows $\forall_F y$ in at $x. \forall t \in T. f\ y\ t < b$
 <proof>

lemma

order-uniform-limitD2:
fixes $l::'a::\text{topological-space} \Rightarrow \text{real}$ — TODO: generalize!
assumes *ul*: *uniform-limit* T f l (at x)
assumes *cont*: *continuous-on* T l
assumes *compact*: *compact* T
assumes *less*: $\bigwedge t. t \in T \implies l\ t > b$
shows $\forall_F y$ in at $x. \forall t \in T. f\ y\ t > b$
 <proof>

lemma *continuous-on-avoid-cases*:

fixes $l::'b::\text{topological-space} \Rightarrow 'a::\text{linear-continuum-topology}$ — TODO: generalize!
assumes *cont*: *continuous-on* T l **and** *conn*: *connected* T
assumes *avoid*: $\bigwedge t. t \in T \implies l\ t \neq b$
obtains $\bigwedge t. t \in T \implies l\ t < b \mid \bigwedge t. t \in T \implies l\ t > b$
 <proof>

lemma

order-uniform-limit-ne:
fixes $l::'a::\text{topological-space} \Rightarrow \text{real}$ — TODO: generalize!
assumes *ul*: *uniform-limit* T f l (at x)
assumes *cont*: *continuous-on* T l
assumes *compact*: *compact* T **and** *conn*: *connected* T
assumes *ne*: $\bigwedge t. t \in T \implies l\ t \neq b$
shows $\forall_F y$ in at $x. \forall t \in T. f\ y\ t \neq b$
 <proof>

lemma *open-cballE*:

assumes *open* S $x \in S$
obtains e **where** $e > 0$ *cball* x $e \subseteq S$
 <proof>

lemma *pos-half-less*: **fixes** $x::\text{real}$ **shows** $x > 0 \implies x / 2 < x$
 <proof>

lemma *closed-levelset*: *closed* $\{x. s\ x = (c::'a::t1\text{-space})\}$ **if** *continuous-on* $UNIV$ s
 <proof>

lemma *closed-levelset-within*: *closed* $\{x \in S. s\ x = (c::'a::t1\text{-space})\}$ **if** *continuous-on* S s *closed* S
 <proof>

context *c1-on-open-euclidean*
begin

lemma *open-existence-ivlE*:

assumes $t \in \text{existence-ivl0 } x \ t \geq 0$

obtains e **where** $e > 0$ $\text{cball } x \ e \times \{0 \ .. \ t + e\} \subseteq \text{Sigma } X \ \text{existence-ivl0}$
 $\langle \text{proof} \rangle$

lemmas [*derivative-intros*] = *flow0-comp-has-derivative*

lemma *flow-isCont-state-space-comp*[*continuous-intros*]:

$t \ x \in \text{existence-ivl0 } (s \ x) \implies \text{isCont } s \ x \implies \text{isCont } t \ x \implies \text{isCont } (\lambda x. \text{flow0}$
 $(s \ x) \ (t \ x)) \ x$
 $\langle \text{proof} \rangle$

lemma *closed-plane*[*simp*]: *closed* $\{x. x \cdot i = c\}$
 $\langle \text{proof} \rangle$

lemma *flow-tendsto-compose*[*tendsto-intros*]:

assumes $(x \longrightarrow xs) \ F \ (t \longrightarrow ts) \ F$

assumes $ts \in \text{existence-ivl0 } xs$

shows $((\lambda s. \text{flow0 } (x \ s) \ (t \ s)) \longrightarrow \text{flow0 } xs \ ts) \ F$
 $\langle \text{proof} \rangle$

lemma *returns-to-implicit-function*:

fixes $s :: 'a :: \text{euclidean-space} \Rightarrow \text{real}$

assumes $rt: \text{returns-to } \{x \in S. s \ x = 0\} \ x \ (\text{is returns-to } ?P \ x)$

assumes $cS: \text{closed } S$

assumes $Ds: \bigwedge x. (s \ \text{has-derivative } \text{blinfun-apply } (Ds \ x)) \ (at \ x)$

assumes $DsC: \text{isCont } Ds \ (\text{poincare-map } ?P \ x)$

assumes $nz: Ds \ (\text{poincare-map } ?P \ x) \ (f \ (\text{poincare-map } ?P \ x)) \neq 0$

obtains $u \ e$

where $s \ (\text{flow0 } x \ (u \ x)) = 0$

$u \ x = \text{return-time } ?P \ x$

$(\bigwedge y. y \in \text{cball } x \ e \implies s \ (\text{flow0 } y \ (u \ y)) = 0)$

$\text{continuous-on } (\text{cball } x \ e) \ u$

$(\lambda t. (t, u \ t)) \ ' \ \text{cball } x \ e \subseteq \text{Sigma } X \ \text{existence-ivl0}$

$0 < e \ (u \ \text{has-derivative } (- \ \text{blinfun-scaleR-left}$

$(\text{inverse } (\text{blinfun-apply } (Ds \ (\text{poincare-map } ?P \ x)) \ (f \ (\text{poincare-map}$
 $?P \ x)))) \ o_L$

$(Ds \ (\text{poincare-map } ?P \ x) \ o_L \ \text{flowderiv } x \ (\text{return-time } ?P \ x)) \ o_L$
 $\text{embed1-blinfun})) \ (at \ x)$

$\langle \text{proof} \rangle$

lemma (*in auto-ll-on-open*) *f-tendsto*[*tendsto-intros*]:

assumes $g1: (g1 \longrightarrow b1) \ (at \ s \ \text{within } S) \ \text{and } b1 \in X$

shows $((\lambda x. f \ (g1 \ x)) \longrightarrow f \ b1) \ (at \ s \ \text{within } S)$

$\langle \text{proof} \rangle$

lemma *flow-avoids-surface-eventually-at-right-pos:*
assumes $s\ x > 0 \vee s\ x = 0 \wedge \text{blinfun-apply } (Ds\ x)\ (f\ x) > 0$
assumes $x: x \in X$
assumes $Ds: \bigwedge x. (s\ \text{has-derivative } Ds\ x)\ (at\ x)$
assumes $DsC: \bigwedge x. \text{isCont } Ds\ x$
shows $\forall_F t\ \text{in } at\text{-right } 0. s\ (\text{flow0 } x\ t) > (0::real)$
<proof>

lemma *flow-avoids-surface-eventually-at-right-neg:*
assumes $s\ x < 0 \vee s\ x = 0 \wedge \text{blinfun-apply } (Ds\ x)\ (f\ x) < 0$
assumes $x: x \in X$
assumes $Ds: \bigwedge x. (s\ \text{has-derivative } Ds\ x)\ (at\ x)$
assumes $DsC: \bigwedge x. \text{isCont } Ds\ x$
shows $\forall_F t\ \text{in } at\text{-right } 0. s\ (\text{flow0 } x\ t) < (0::real)$
<proof>

lemma *flow-avoids-surface-eventually-at-right:*
assumes $x \notin S \vee s\ x \neq 0 \vee \text{blinfun-apply } (Ds\ x)\ (f\ x) \neq 0$
assumes $x: x \in X$ **and** $cS: \text{closed } S$
assumes $Ds: \bigwedge x. (s\ \text{has-derivative } Ds\ x)\ (at\ x)$
assumes $DsC: \bigwedge x. \text{isCont } Ds\ x$
shows $\forall_F t\ \text{in } at\text{-right } 0. (\text{flow0 } x\ t) \notin \{x \in S. s\ x = (0::real)\}$
<proof>

lemma *eventually-returns-to:*
fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes $rt: \text{returns-to } \{x \in S. s\ x = 0\}\ x\ (\text{is returns-to } ?P\ x)$
assumes $cS: \text{closed } S$
assumes $Ds: \bigwedge x. (s\ \text{has-derivative } \text{blinfun-apply } (Ds\ x))\ (at\ x)$
assumes $DsC: \bigwedge x. \text{isCont } Ds\ x$
assumes $\text{eventually-inside: } \forall_F x\ \text{in } at\ (\text{poincare-map } ?P\ x). s\ x = 0 \longrightarrow x \in S$
assumes $nz: Ds\ (\text{poincare-map } ?P\ x)\ (f\ (\text{poincare-map } ?P\ x)) \neq 0$
assumes $nz0: x \notin S \vee s\ x \neq 0 \vee Ds\ x\ (f\ x) \neq 0$
shows $\forall_F x\ \text{in } at\ x. \text{returns-to } ?P\ x$
<proof>

lemma
return-time-isCont-outside:
fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes $rt: \text{returns-to } \{x \in S. s\ x = 0\}\ x\ (\text{is returns-to } ?P\ x)$
assumes $cS: \text{closed } S$
assumes $Ds: \bigwedge x. (s\ \text{has-derivative } \text{blinfun-apply } (Ds\ x))\ (at\ x)$
assumes $DsC: \bigwedge x. \text{isCont } Ds\ x$
assumes $\text{through: } (Ds\ (\text{poincare-map } ?P\ x))\ (f\ (\text{poincare-map } ?P\ x)) \neq 0$
assumes $\text{eventually-inside: } \forall_F x\ \text{in } at\ (\text{poincare-map } ?P\ x). s\ x = 0 \longrightarrow x \in S$
assumes $\text{outside: } x \notin S \vee s\ x \neq 0$
shows $\text{isCont } (\text{return-time } ?P)\ x$
<proof>

lemma *isCont-poincare-map*:
assumes *isCont* (return-time P) x
returns-to P x closed P
shows *isCont* (poincare-map P) x
 \langle proof \rangle

lemma *poincare-map-tendsto*:
assumes (return-time $P \longrightarrow$ return-time P x) (at x within S)
returns-to P x closed P
shows (poincare-map $P \longrightarrow$ poincare-map P x) (at x within S)
 \langle proof \rangle

lemma
return-time-continuous-below:
fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes *rt*: *returns-to* $\{x \in S. s\ x = 0\}$ x (**is** *returns-to* $?P$ x)
assumes *Ds*: $\bigwedge x. (s \text{ has-derivative } \text{blinfun-apply } (Ds\ x))$ (at x)
assumes *cS*: closed S
assumes *eventually-inside*: $\forall_F x$ in at (poincare-map $?P$ x). $s\ x = 0 \longrightarrow x \in S$
assumes *DsC*: $\bigwedge x. \text{isCont } Ds\ x$
assumes *through*: $(Ds\ (\text{poincare-map } ?P\ x)) (f\ (\text{poincare-map } ?P\ x)) \neq 0$
assumes *inside*: $x \in S\ s\ x = 0\ Ds\ x\ (f\ x) < 0$
shows *continuous* (at x within $\{x. s\ x \leq 0\}$) (return-time $?P$)
 \langle proof \rangle

lemma
return-time-continuous-below-plane:
fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes *rt*: *returns-to* $\{x \in R. x \cdot n = c\}$ x (**is** *returns-to* $?P$ x)
assumes *cR*: closed R
assumes *through*: $f\ (\text{poincare-map } ?P\ x) \cdot n \neq 0$
assumes *R*: $x \in R$
assumes *inside*: $x \cdot n = c\ f\ x \cdot n < 0$
assumes *eventually-inside*: $\forall_F x$ in at (poincare-map $?P$ x). $x \cdot n = c \longrightarrow x \in R$
shows *continuous* (at x within $\{x. x \cdot n \leq c\}$) (return-time $?P$)
 \langle proof \rangle

lemma
poincare-map-in-interior-eventually-return-time-equal:
assumes *RP*: $R \subseteq P$
assumes *cP*: closed P
assumes *cR*: closed R
assumes *ret*: *returns-to* P x
assumes *evret*: $\forall_F x$ in at x within S . *returns-to* P x
assumes *evR*: $\forall_F x$ in at x within S . *poincare-map* P $x \in R$
shows $\forall_F x$ in at x within S . *returns-to* R $x \wedge$ return-time P $x =$ return-time R x
 \langle proof \rangle

lemma *poincare-map-in-planeI*:

assumes *returns-to* (plane n c) $x0$
shows *poincare-map* (plane n c) $x0 \cdot n = c$
 ⟨*proof*⟩

lemma *less-return-time-imp-exivl*:

$h \in \text{existence-ivl0 } x'$ **if** $h \leq \text{return-time } P x'$ *returns-to* $P x'$ *closed* $P 0 \leq h$
 ⟨*proof*⟩

lemma *eventually-returns-to-continuousI*:

assumes *returns-to* $P x$
assumes *closed* P
assumes *continuous* (at x within S) (return-time P)
shows $\forall_F x$ in at x within S . *returns-to* $P x$
 ⟨*proof*⟩

lemma *return-time-implicit-functionE*:

fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes *rt*: *returns-to* $\{x \in S. s x = 0\} x$ (**is** *returns-to* $?P -$)
assumes *cS*: *closed* S
assumes *Ds*: $\bigwedge x. (s \text{ has-derivative } \text{blinfun-apply } (Ds x)) (at x)$
assumes *DsC*: $\bigwedge x. \text{isCont } Ds x$
assumes *Ds-through*: $(Ds (\text{poincare-map } ?P x)) (f (\text{poincare-map } ?P x)) \neq 0$
assumes *eventually-inside*: $\forall_F x$ in at (poincare-map $?P x$). $s x = 0 \longrightarrow x \in S$
assumes *outside*: $x \notin S \vee s x \neq 0$
obtains e' **where**
 $0 < e'$
 $\bigwedge y. y \in \text{ball } x e' \implies \text{returns-to } ?P y$
 $\bigwedge y. y \in \text{ball } x e' \implies s (\text{flow0 } y (\text{return-time } ?P y)) = 0$
continuous-on (ball $x e'$) (return-time $?P$)
 $(\bigwedge y. y \in \text{ball } x e' \implies Ds (\text{poincare-map } ?P y) \text{ o}_L \text{flowderiv } y (\text{return-time } ?P y)) \text{ o}_L \text{embed2-blinfun} \in \text{invertibles-blinfun}$
 $(\bigwedge U v sa. (\bigwedge sa. sa \in U \implies s (\text{flow0 } sa (v sa)) = 0) \implies \text{return-time } ?P x = v x \implies \text{continuous-on } U v \implies sa \in U \implies x \in U \implies U \subseteq \text{ball } x e' \implies \text{connected } U \implies \text{open } U \implies \text{return-time } ?P sa = v sa)$
 (return-time $?P$ has-derivative
 – *blinfun-scaleR-left* (inverse $((Ds (\text{poincare-map } ?P x)) (f (\text{poincare-map } ?P x)))) \text{ o}_L (Ds (\text{poincare-map } ?P x) \text{ o}_L D\text{flow } x (\text{return-time } ?P x))$)
 (at x)
 ⟨*proof*⟩

lemma *return-time-has-derivative*:

fixes $s::'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes *rt*: *returns-to* $\{x \in S. s x = 0\} x$ (**is** *returns-to* $?P -$)
assumes *cS*: *closed* S

assumes Ds : $\bigwedge x. (s \text{ has-derivative } \text{blinfun-apply } (Ds \ x)) \ (at \ x)$
assumes DsC : $\bigwedge x. \text{isCont } Ds \ x$
assumes $Ds\text{-through}$: $(Ds \ (\text{poincare-map } ?P \ x)) \ (f \ (\text{poincare-map } ?P \ x)) \neq 0$
assumes eventually-inside : $\forall_F \ x \ \text{in } at \ (\text{poincare-map } \{x \in S. \ s \ x = 0\} \ x). \ s \ x = 0 \longrightarrow x \in S$
assumes outside : $x \notin S \vee s \ x \neq 0$
shows $(\text{return-time } ?P \ \text{has-derivative}$
 $\quad - \text{blinfun-scaleR-left } (\text{inverse } ((Ds \ (\text{poincare-map } ?P \ x)) \ (f \ (\text{poincare-map } ?P \ x)))) \ o_L$
 $\quad (Ds \ (\text{poincare-map } ?P \ x) \ o_L \ Dflow \ x \ (\text{return-time } ?P \ x)))$
 $\quad (at \ x)$
 $\langle \text{proof} \rangle$

lemma $\text{return-time-plane-has-derivative-blinfun}$:

assumes rt : $\text{returns-to } \{x \in S. \ x \cdot i = c\} \ x \ (\text{is } \text{returns-to } ?P \ -)$
assumes cS : $\text{closed } S$
assumes fnz : $f \ (\text{poincare-map } ?P \ x) \cdot i \neq 0$
assumes eventually-inside : $\forall_F \ x \ \text{in } at \ (\text{poincare-map } ?P \ x). \ x \cdot i = c \longrightarrow x \in S$
assumes outside : $x \notin S \vee x \cdot i \neq c$
shows $(\text{return-time } ?P \ \text{has-derivative}$
 $\quad - \text{blinfun-scaleR-left } (\text{inverse } ((\text{blinfun-inner-left } i) \ (f \ (\text{poincare-map } ?P \ x))))$
 $\quad o_L$
 $\quad (\text{blinfun-inner-left } i \ o_L \ Dflow \ x \ (\text{return-time } ?P \ x))) \ (at \ x)$
 $\langle \text{proof} \rangle$

lemma $\text{return-time-plane-has-derivative}$:

assumes rt : $\text{returns-to } \{x \in S. \ x \cdot i = c\} \ x \ (\text{is } \text{returns-to } ?P \ -)$
assumes cS : $\text{closed } S$
assumes fnz : $f \ (\text{poincare-map } ?P \ x) \cdot i \neq 0$
assumes eventually-inside : $\forall_F \ x \ \text{in } at \ (\text{poincare-map } ?P \ x). \ x \cdot i = c \longrightarrow x \in S$
assumes outside : $x \notin S \vee x \cdot i \neq c$
shows $(\text{return-time } ?P \ \text{has-derivative}$
 $\quad (\lambda h. \ - \ (Dflow \ x \ (\text{return-time } ?P \ x)) \ h \cdot i / (f \ (\text{poincare-map } ?P \ x) \cdot i))) \ (at \ x)$
 $\langle \text{proof} \rangle$

definition $D\text{poincare-map } i \ c \ S \ x =$

$(\lambda h. \ (Dflow \ x \ (\text{return-time } \{x \in S. \ x \cdot i = c\} \ x)) \ h \ -$
 $\quad ((Dflow \ x \ (\text{return-time } \{x \in S. \ x \cdot i = c\} \ x)) \ h \cdot i /$
 $\quad (f \ (\text{poincare-map } \{x \in S. \ x \cdot i = c\} \ x) \cdot i)) \ *_R \ f \ (\text{poincare-map } \{x \in S. \ x$
 $\cdot i = c\} \ x))$

definition $D\text{poincare-map}' \ i \ c \ S \ x =$

$Dflow \ x \ (\text{return-time } \{x \in S. \ x \cdot i - c = 0\} \ x) \ -$
 $(\text{blinfun-scaleR-left } (f \ (\text{poincare-map } \{x \in S. \ x \cdot i = c\} \ x)) \ o_L$
 $\quad (\text{blinfun-scaleR-left } (\text{inverse } ((f \ (\text{poincare-map } \{x \in S. \ x \cdot i = c\} \ x) \cdot i))) \ o_L$
 $\quad (\text{blinfun-inner-left } i \ o_L \ Dflow \ x \ (\text{return-time } \{x \in S. \ x \cdot i - c = 0\} \ x))))$

theorem $\text{poincare-map-plane-has-derivative}$:

assumes rt : $\text{returns-to } \{x \in S. \ x \cdot i = c\} \ x \ (\text{is } \text{returns-to } ?P \ -)$

assumes cS : *closed S*
assumes fnz : f (*poincare-map ?P x*) $\cdot i \neq 0$
assumes *eventually-inside*: $\forall_F x$ *in at* (*poincare-map ?P x*). $x \cdot i = c \longrightarrow x \in S$
assumes *outside*: $x \notin S \vee x \cdot i \neq c$
notes [*derivative-intros*] = *return-time-plane-has-derivative[OF rt cS fnz eventually-inside outside]*
shows (*poincare-map ?P has-derivative Dpoincare-map' i c S x*) (*at x*)
<proof>

end

end

theory *Reachability-Analysis*

imports

Flow

Poincare-Map

begin

lemma *not-mem-eq-mem-not*: $a \notin A \longleftrightarrow a \in - A$
<proof>

lemma *continuous-orderD*:
fixes $g :: 'b :: t2-space \Rightarrow 'c :: order-topology$
assumes *continuous* (*at x within S*) g
shows $g x > c \Longrightarrow \forall_F y$ *in at x within S*. $g y > c$
 $g x < c \Longrightarrow \forall_F y$ *in at x within S*. $g y < c$
<proof>

lemma *frontier-halfspace-component-ge*: $n \neq 0 \Longrightarrow \text{frontier } \{x. c \leq x \cdot n\} = \text{plane } n c$
<proof>

lemma *closed-Collect-le-within*:
fixes $f g :: 'a :: topological-space \Rightarrow 'b :: linorder-topology$
assumes f : *continuous-on UNIV* f
and g : *continuous-on UNIV* g
and *closed* R
shows *closed* $\{x \in R. f x \leq g x\}$
<proof>

6.1 explicit representation of hyperplanes / halfspaces

datatype $'a$ *sctn* = *Sctn* (*normal*: $'a$) (*pstn*: *real*)

definition *le-halfspace sctn* $x \longleftrightarrow x \cdot \text{normal } sctn \leq \text{pstn } sctn$

definition *lt-halfspace sctn* $x \longleftrightarrow x \cdot \text{normal } sctn < \text{pstn } sctn$

definition *ge-halfspace sctn* $x \longleftrightarrow x \cdot \text{normal } sctn \geq \text{pstn } sctn$

definition *gt-halfspace* *sctn* $x \longleftrightarrow x \cdot \text{normal } sctn > \text{pstrn } sctn$

definition *plane-of* *sctn* = $\{x. x \cdot \text{normal } sctn = \text{pstrn } sctn\}$

definition *above-halfspace* *sctn* = *Collect* (*ge-halfspace* *sctn*)

definition *below-halfspace* *sctn* = *Collect* (*le-halfspace* *sctn*)

definition *sbelow-halfspace* *sctn* = *Collect* (*lt-halfspace* *sctn*)

definition *sabove-halfspace* *sctn* = *Collect* (*gt-halfspace* *sctn*)

6.2 explicit H representation of polytopes (mind *Polytopes.thy*)

definition *below-halfspaces*

where *below-halfspaces* *sctns* = $\bigcap (\text{below-halfspace } 'sctns)$

definition *sbelow-halfspaces*

where *sbelow-halfspaces* *sctns* = $\bigcap (\text{sbelow-halfspace } 'sctns)$

definition *above-halfspaces*

where *above-halfspaces* *sctns* = $\bigcap (\text{above-halfspace } 'sctns)$

definition *sabove-halfspaces*

where *sabove-halfspaces* *sctns* = $\bigcap (\text{sabove-halfspace } 'sctns)$

lemmas *halfspace-simps* =

above-halfspace-def

sabove-halfspace-def

below-halfspace-def

sbelow-halfspace-def

below-halfspaces-def

sbelow-halfspaces-def

above-halfspaces-def

sabove-halfspaces-def

ge-halfspace-def[*abs-def*]

gt-halfspace-def[*abs-def*]

le-halfspace-def[*abs-def*]

lt-halfspace-def[*abs-def*]

6.3 predicates for reachability analysis

context *c1-on-open-euclidean*

begin

definition *flowpipe* ::

$((a::\text{euclidean-space}) \times (a \Rightarrow_L a)) \text{ set} \Rightarrow \text{real} \Rightarrow \text{real} \Rightarrow$
 $(a \times (a \Rightarrow_L a)) \text{ set} \Rightarrow (a \times (a \Rightarrow_L a)) \text{ set} \Rightarrow \text{bool}$

where *flowpipe* $X0\ hl\ hu\ CX\ X1 \iff 0 \leq hl \wedge hl \leq hu \wedge fst\ 'X0 \subseteq X \wedge fst\ 'CX \subseteq X \wedge fst\ 'X1 \subseteq X \wedge$
 $(\forall (x0, d0) \in X0. \forall h \in \{hl .. hu\}.$
 $h \in existence-ivl0\ x0 \wedge (flow0\ x0\ h, Dflow\ x0\ h\ o_L\ d0) \in X1 \wedge (\forall h' \in \{0 .. h\}.$
 $(flow0\ x0\ h', Dflow\ x0\ h'\ o_L\ d0) \in CX))$

lemma *flowpipeD*:

assumes *flowpipe* $X0\ hl\ hu\ CX\ X1$

shows *flowpipe-safeD*: $fst\ 'X0 \cup fst\ 'CX \cup fst\ 'X1 \subseteq X$

and *flowpipe-nonneg*: $0 \leq hl\ hl \leq hu$

and *flowpipe-exivl*: $hl \leq h \implies h \leq hu \implies (x0, d0) \in X0 \implies h \in existence-ivl0\ x0$

and *flowpipe-discrete*: $hl \leq h \implies h \leq hu \implies (x0, d0) \in X0 \implies (flow0\ x0\ h,$
 $Dflow\ x0\ h\ o_L\ d0) \in X1$

and *flowpipe-cont*: $hl \leq h \implies h \leq hu \implies (x0, d0) \in X0 \implies 0 \leq h' \implies h' \leq h \implies (flow0\ x0\ h', Dflow\ x0\ h'\ o_L\ d0) \in CX$

<proof>

lemma *flowpipe-source-subset*: *flowpipe* $X0\ hl\ hu\ CX\ X1 \implies X0 \subseteq CX$

<proof>

definition *flowsto* $X0\ T\ CX\ X1 \iff$

$(\forall (x0, d0) \in X0. \exists h \in T. h \in existence-ivl0\ x0 \wedge (flow0\ x0\ h, Dflow\ x0\ h\ o_L\ d0) \in X1 \wedge (\forall h' \in open-segment\ 0\ h. (flow0\ x0\ h', Dflow\ x0\ h'\ o_L\ d0) \in CX))$

lemma *flowsto-to-empty-iff[simp]*: *flowsto* $a\ t\ b\ \{\}$ $\iff a = \{\}$

<proof>

lemma *flowsto-from-empty-iff[simp]*: *flowsto* $\{\}\ t\ b\ c$

<proof>

lemma *flowsto-empty-time-iff[simp]*: *flowsto* $a\ \{\}\ b\ c \iff a = \{\}$

<proof>

lemma *flowstoE*:

assumes *flowsto* $X0\ T\ CX\ X1\ (x0, d0) \in X0$

obtains h **where** $h \in T\ h \in existence-ivl0\ x0\ (flow0\ x0\ h, Dflow\ x0\ h\ o_L\ d0) \in X1$

$\wedge h'. h' \in open-segment\ 0\ h \implies (flow0\ x0\ h', Dflow\ x0\ h'\ o_L\ d0) \in CX$

<proof>

lemma *flowsto-safeD*: *flowsto* $X0\ T\ CX\ X1 \implies fst\ 'X0 \subseteq X$

<proof>

lemma *flowsto-union*:

assumes *1*: *flowsto* $X0\ T\ CX\ Y$ **and** *2*: *flowsto* $Z\ S\ CZ\ W$

shows *flowsto* $(X0 \cup Z)\ (T \cup S)\ (CX \cup CZ)\ (Y \cup W)$

<proof>

lemma *flowsto-subset*:

assumes *flowsto* $X0\ T\ CX\ Y$

assumes $Z \subseteq X0\ T \subseteq S\ CX \subseteq CZ\ Y \subseteq W$

shows *flowsto* $Z\ S\ CZ\ W$

<proof>

lemmas *flowsto-unionI* = *flowsto-subset*[*OF flowsto-union*]

lemma *flowsto-unionE*:

assumes *flowsto* $X0\ T\ CX\ (Y \cup Z)$

obtains $X1\ X2$ **where** $X0 = X1 \cup X2$ *flowsto* $X1\ T\ CX\ Y$ *flowsto* $X2\ T\ CX\ Z$

<proof>

lemma *flowsto-trans*:

assumes A : *flowsto* $A\ S\ B\ C$ **and** C : *flowsto* $C\ T\ D\ E$

shows *flowsto* $A\ \{s + t \mid s \in S \wedge t \in T\}\ (B \cup D \cup C)\ E$

<proof>

lemma *flowsto-step*:

assumes A : *flowsto* $A\ S\ B\ C$

assumes D : *flowsto* $D\ T\ E\ F$

shows *flowsto* $A\ (S \cup \{s + t \mid s \in S \wedge t \in T\})\ (B \cup E \cup C \cap D)\ (C - D \cup F)$

<proof>

lemma

flowsto-stepI:

flowsto $X0\ U\ B\ C \implies$

flowsto $D\ T\ E\ F \implies$

$Z \subseteq X0 \implies$

$(\bigwedge s. s \in U \implies s \in S) \implies$

$(\bigwedge s\ t. s \in U \implies t \in T \implies s + t \in S) \implies$

$B \cup E \cup D \cap C \subseteq CZ \implies C - D \cup F \subseteq W \implies$ *flowsto* $Z\ S\ CZ\ W$

<proof>

lemma *flowsto-imp-flowsto*:

flowpipe $Y\ h\ h\ CY\ Z \implies$ *flowsto* $Y\ \{h\}\ (CY)\ Z$

<proof>

lemma *connected-below-halfspace*:

assumes $x \in$ *below-halfspace* $sctn$

assumes $x \in S$ *connected* S

assumes $S \cap$ *plane-of* $sctn = \{\}$

shows $S \subseteq$ *below-halfspace* $sctn$

<proof>

lemma

inter-Collect-eq-empty:

assumes $\bigwedge x. x \in X0 \implies \neg g\ x$ **shows** $X0 \cap$ *Collect* $g = \{\}$

<proof>

6.4 Poincare Map

lemma *closed-plane-of[simp]*: *closed (plane-of sctn)*

<proof>

definition *poincare-mapsto* $P X0 S CX Y \longleftrightarrow (\forall (x, d) \in X0.$

returns-to $P x \wedge \text{fst } 'X0 \subseteq S \wedge$

(return-time P *differentiable at* x *within* $S) \wedge$

$(\exists D. (\text{poincare-map } P \text{ has-derivative } \text{blinfun-apply } D) (\text{at } x \text{ within } S) \wedge$

$(\text{poincare-map } P x, D \text{ o}_L d) \in Y) \wedge$

$(\forall t \in \{0 <.. < \text{return-time } P x\}. \text{flow0 } x t \in CX))$

lemma *poincare-mapsto-empty[simp]*:

poincare-mapsto $P \{\} S CX Y$

<proof>

lemma *flowsto-eventually-mem-cont*:

assumes *flowsto* $X0 T CX Y (x, d) \in X0 T \subseteq \{0 <..\}$

shows $\forall_F t$ *in at-right* $0. (\text{flow0 } x t, D\text{flow } x t \text{ o}_L d) \in CX$

<proof>

lemma *frontier-aux-lemma*:

fixes $R :: 'n::\text{euclidean-space set}$

assumes *closed* $R R \subseteq \{x. x \cdot n = c\}$ **and** *[simp]*: $n \neq 0$

shows *frontier* $\{x \in R. c \leq x \cdot n\} = \{x \in R. c = x \cdot n\}$

<proof>

lemma *blinfun-minus-comp-distrib*: $(a - b) \text{ o}_L c = (a \text{ o}_L c) - (b \text{ o}_L c)$

<proof>

lemma *flowpipe-split-at-above-halfspace*:

assumes *flowpipe* $X0 \text{ hl } t CX Y \text{fst } 'X0 \cap \{x. x \cdot n \geq c\} = \{\}$ **and** *[simp]*: $n \neq 0$

assumes *cR*: *closed* R **and** *Rs*: $R \subseteq \text{plane } n c$

assumes *PDP*: $\bigwedge x d. (x, d) \in CX \implies x \cdot n = c \implies (x,$

$d - (\text{blinfun-scaleR-left } (f (x)) \text{ o}_L (\text{blinfun-scaleR-left } (\text{inverse } (f x \cdot n)) \text{ o}_L$

$(\text{blinfun-inner-left } n \text{ o}_L d))) \in \text{PDP}$

assumes *PDP-nz*: $\bigwedge x d. (x, d) \in \text{PDP} \implies f x \cdot n \neq 0$

assumes *PDP-inR*: $\bigwedge x d. (x, d) \in \text{PDP} \implies x \in R$

assumes *PDP-in*: $\bigwedge x d. (x, d) \in \text{PDP} \implies \forall_F x$ *in at* x *within plane* $n c. x \in R$

obtains $X1 X2$ **where** $X0 = X1 \cup X2$

flowsto $X1 \{0 <..t\} (CX \cap \{x. x \cdot n < c\} \times \text{UNIV}) (CX \cap \{x \in R. x \cdot n =$

$c\} \times \text{UNIV})$

flowsto $X2 \{\text{hl } .. t\} (CX \cap \{x. x \cdot n < c\} \times \text{UNIV}) (Y \cap (\{x. x \cdot n < c\} \times$

$\text{UNIV}))$

poincare-mapsto $\{x \in R. x \cdot n = c\} X1 \text{UNIV } (\text{fst } 'CX \cap \{x. x \cdot n < c\}) \text{PDP}$

<proof>

lemma *poincare-map-has-derivative-step:*

assumes *Deriv:* (*poincare-map P has-derivative blinfun-apply D*) (*at (flow0 x0 h)*)

assumes *ret:* *returns-to P x0*

assumes *cont:* *continuous (at x0 within S) (return-time P)*

assumes *less:* $0 \leq h \wedge h < \text{return-time } P \ x0$

assumes *cP:* *closed P and x0: x0 ∈ S*

shows ($\lambda x. \text{poincare-map } P \ x$) *has-derivative (D o_L Dflow x0 h)* (*at x0 within S*)

<proof>

lemma *poincare-mapsto-trans:*

assumes *poincare-mapsto p1 X0 S CX P1*

assumes *poincare-mapsto p2 P1 UNIV CY P2*

assumes $CX \cup CY \cup \text{fst } ' P1 \subseteq CZ$

assumes $p2 \cap (CX \cup \text{fst } ' P1) = \{\}$

assumes [*intro, simp*]: *closed p1*

assumes [*intro, simp*]: *closed p2*

assumes *cont:* $\bigwedge x \ d. (x, d) \in X0 \implies \text{continuous (at } x \text{ within } S) \text{ (return-time } p2)$

shows *poincare-mapsto p2 X0 S CZ P2*

<proof>

lemma *flowsto-poincare-trans:*— **TODO:** the proof is close to $\llbracket \text{poincare-mapsto } ?p1.0 \ ?X0.0 \ ?S \ ?CX \ ?P1.0; \text{poincare-mapsto } ?p2.0 \ ?P1.0 \ \text{UNIV } ?CY \ ?P2.0; \ ?CX \cup \ ?CY \cup \text{fst } ' \ ?P1.0 \subseteq \ ?CZ; \ ?p2.0 \cap (\ ?CX \cup \text{fst } ' \ ?P1.0) = \{\}; \text{closed } ?p1.0; \text{closed } ?p2.0; \bigwedge x \ d. (x, d) \in \ ?X0.0 \implies \text{continuous (at } x \text{ within } ?S) \text{ (return-time } ?p2.0) \rrbracket \implies \text{poincare-mapsto } ?p2.0 \ ?X0.0 \ ?S \ ?CZ \ ?P2.0$

assumes *f:* *flowsto X0 T CX P1*

assumes *poincare-mapsto p2 P1 UNIV CY P2*

assumes *nn:* $\bigwedge t. t \in T \implies t \geq 0$

assumes $\text{fst } ' CX \cup CY \cup \text{fst } ' P1 \subseteq CZ$

assumes $p2 \cap (\text{fst } ' CX \cup \text{fst } ' P1) = \{\}$

assumes [*intro, simp*]: *closed p2*

assumes *cont:* $\bigwedge x \ d. (x, d) \in X0 \implies \text{continuous (at } x \text{ within } S) \text{ (return-time } p2)$

assumes *subset:* $\text{fst } ' X0 \subseteq S$

shows *poincare-mapsto p2 X0 S CZ P2*

<proof>

6.5 conditions for continuous return time

definition *section s Ds S* \longleftrightarrow

$(\forall x. (s \text{ has-derivative blinfun-apply } (Ds \ x)) \text{ (at } x)) \wedge$

$(\forall x. \text{isCont } Ds \ x) \wedge$

$(\forall x \in S. s \ x = (0::\text{real}) \longrightarrow Ds \ x (f \ x) \neq 0) \wedge$

$\text{closed } S \wedge S \subseteq X$

lemma *sectionD*:

assumes *section s Ds S*

shows (*s has-derivative blinfun-apply (Ds x)*) (*at x*)

isCont Ds x

$x \in S \implies s\ x = 0 \implies Ds\ x\ (f\ x) \neq 0$

closed S S $\subseteq X$

<proof>

definition *transversal p* $\longleftrightarrow (\forall x \in p. \forall_F t\ \text{in}\ \text{at-right}\ 0. \text{flow0}\ x\ t \notin p)$

lemma *transversalD*: *transversal p* $\implies x \in p \implies \forall_F t\ \text{in}\ \text{at-right}\ 0. \text{flow0}\ x\ t \notin p$

<proof>

lemma *transversal-section*:

fixes *c::real*

assumes *section s Ds S*

shows *transversal* $\{x \in S. s\ x = 0\}$

<proof>

lemma *section-closed[intro, simp]*: *section s Ds S* $\implies \text{closed}\ \{x \in S. s\ x = 0\}$

<proof>

lemma *return-time-continuous-belowI*:

assumes *ft: flowsto X0 T CX X1*

assumes *pos*: $\bigwedge t. t \in T \implies t > 0$

assumes *X0*: *fst* ' *X0* $\subseteq \{x \in S. s\ x = 0\}$

assumes *CX*: *fst* ' *CX* $\cap \{x \in S. s\ x = 0\} = \{\}$

assumes *X1*: *fst* ' *X1* $\subseteq \{x \in S. s\ x = 0\}$

assumes *sec*: *section s Ds S*

assumes *nz*: $\bigwedge x. x \in S \implies s\ x = 0 \implies Ds\ x\ (f\ x) \neq 0$

assumes *Dneg*: $(\lambda x. (Ds\ x)\ (f\ x))$ ' *X0* $\subseteq \{..<0\}$

assumes *rel-int*: $\bigwedge x. x \in \text{fst}'\ X1 \implies \forall_F x\ \text{in}\ \text{at}\ x. s\ x = 0 \longrightarrow x \in S$

assumes $(x, d) \in X0$

shows *continuous* (*at x within* $\{x. s\ x \leq 0\}$) (*return-time* $\{x \in S. s\ x = 0\}$)

<proof>

end

end

theory *Flow-Congs*

imports *Reachability-Analysis*

begin

lemma *lipschitz-on-congI*:

assumes *L'-lipschitz-on s' g'*

assumes $s' = s$

assumes $L' \leq L$
assumes $\bigwedge x y. x \in s \implies g' x = g x$
shows $L\text{-lipschitz-on } s g$
 $\langle \text{proof} \rangle$

lemma *local-lipschitz-congI*:
assumes *local-lipschitz* $s' t' g'$
assumes $s' = s$
assumes $t' = t$
assumes $\bigwedge x y. x \in s \implies y \in t \implies g' x y = g x y$
shows *local-lipschitz* $s t g$
 $\langle \text{proof} \rangle$

context *ll-on-open-it*— TODO: do this more generically for *ll-on-open-it*
begin

context **fixes** $S Y g$ **assumes** *cong*: $X = Y T = S \bigwedge x t. x \in Y \implies t \in S \implies f t x = g t x$
begin

lemma *ll-on-open-congI*: *ll-on-open* $S g Y$
 $\langle \text{proof} \rangle$

lemma *existence-ivl-subsetI*:
assumes $t \in \text{existence-ivl } t0 x0$
shows $t \in \text{ll-on-open.existence-ivl } S g Y t0 x0$
 $\langle \text{proof} \rangle$

lemma *existence-ivl-cong*:
shows $\text{existence-ivl } t0 x0 = \text{ll-on-open.existence-ivl } S g Y t0 x0$
 $\langle \text{proof} \rangle$

lemma *flow-cong*:
assumes $t \in \text{existence-ivl } t0 x0$
shows $\text{flow } t0 x0 t = \text{ll-on-open.flow } S g Y t0 x0 t$
 $\langle \text{proof} \rangle$

end

end

context *auto-ll-on-open* **begin**

context **fixes** $Y g$ **assumes** *cong*: $X = Y \bigwedge x t. x \in Y \implies f x = g x$
begin

lemma *auto-ll-on-open-congI*: *auto-ll-on-open* $g Y$
 $\langle \text{proof} \rangle$

lemma *existence-ivl0-cong*:
shows *existence-ivl0* $x0 = \text{auto-ll-on-open.existence-ivl0 } g \ Y \ x0$
 $\langle \text{proof} \rangle$

lemma *flow0-cong*:
assumes $t \in \text{existence-ivl0 } x0$
shows *flow0* $x0 \ t = \text{auto-ll-on-open.flow0 } g \ Y \ x0 \ t$
 $\langle \text{proof} \rangle$

end

end

context *c1-on-open-euclidean* **begin**

context **fixes** $Y \ g$ **assumes** *cong*: $X = Y \ \bigwedge x \ t. \ x \in Y \ \Longrightarrow \ f \ x = g \ x$
begin

lemma *f'-cong*: (*g* has-derivative *blinfun-apply* ($f' \ x$)) (at x) **if** $x \in Y$
 $\langle \text{proof} \rangle$

lemma *c1-on-open-euclidean-congI*: *c1-on-open-euclidean* $g \ f' \ Y$
 $\langle \text{proof} \rangle$

lemma *vareq-cong*: *vareq* $x0 \ t = \text{c1-on-open-euclidean.vareq } g \ f' \ Y \ x0 \ t$
if $t \in \text{existence-ivl0 } x0$
 $\langle \text{proof} \rangle$

lemma *Dflow-cong*:
assumes $t \in \text{existence-ivl0 } x0$
shows *Dflow* $x0 \ t = \text{c1-on-open-euclidean.Dflow } g \ f' \ Y \ x0 \ t$
 $\langle \text{proof} \rangle$

lemma *flowsto-congI1*:
assumes *flowsto* $A \ B \ C \ D$
shows *c1-on-open-euclidean.flowsto* $g \ f' \ Y \ A \ B \ C \ D$
 $\langle \text{proof} \rangle$

lemma *flowsto-congI2*:
assumes *c1-on-open-euclidean.flowsto* $g \ f' \ Y \ A \ B \ C \ D$
shows *flowsto* $A \ B \ C \ D$
 $\langle \text{proof} \rangle$

lemma *flowsto-congI*: *flowsto* $A \ B \ C \ D = \text{c1-on-open-euclidean.flowsto } g \ f' \ Y \ A \ B \ C \ D$
 $\langle \text{proof} \rangle$

lemma

returns-to-congI1:
assumes *returns-to A x*
shows *auto-ll-on-open.returns-to g Y A x*
 ⟨*proof*⟩

lemma
returns-to-congI2:
assumes *auto-ll-on-open.returns-to g Y x A*
shows *returns-to x A*
 ⟨*proof*⟩

lemma *returns-to-cong: auto-ll-on-open.returns-to g Y A x = returns-to A x*
 ⟨*proof*⟩

lemma
return-time-cong:
shows *return-time A x = auto-ll-on-open.return-time g Y A x*
 ⟨*proof*⟩

lemma *poincare-mapsto-congI1:*
assumes *poincare-mapsto A B C D E closed A*
shows *c1-on-open-euclidean.poincare-mapsto g Y A B C D E*
 ⟨*proof*⟩

lemma *poincare-mapsto-congI2:*
assumes *c1-on-open-euclidean.poincare-mapsto g Y A B C D E closed A*
shows *poincare-mapsto A B C D E*
 ⟨*proof*⟩

lemma *poincare-mapsto-cong: closed A \implies*
poincare-mapsto A B C D E = c1-on-open-euclidean.poincare-mapsto g Y A B
C D E
 ⟨*proof*⟩

end

end

end

theory *Cones*

imports

HOL-Analysis.Analysis

Triangle.Triangle

../ODE-Auxiliarities

begin

lemma *arcsin-eq-zero-iff[simp]: $-1 \leq x \implies x \leq 1 \implies \arcsin x = 0 \iff x = 0$*
 ⟨*proof*⟩

definition *conemem* :: 'a::real-vector \Rightarrow 'a \Rightarrow real \Rightarrow 'a **where** *conemem* u v t =
 $\cos t *_R u + \sin t *_R v$

definition *conesegment* u v = *conemem* u v ' {0.. pi / 2}

lemma

bounded-linear-image-conemem:

assumes *bounded-linear* F

shows F (conemem u v t) = conemem (F u) (F v) t

<proof>

lemma

bounded-linear-image-conesegment:

assumes *bounded-linear* F

shows F ' conesegment u v = conesegment (F u) (F v)

<proof>

lemma *discriminant*: $a * x^2 + b * x + c = (0::real) \implies 0 \leq b^2 - 4 * a * c$
<proof>

lemma *quadratic-eq-factoring*:

assumes D: $D = b^2 - 4 * a * c$

assumes nn: $0 \leq D$

assumes x1: $x_1 = (-b + \text{sqr}t D) / (2 * a)$

assumes x2: $x_2 = (-b - \text{sqr}t D) / (2 * a)$

assumes a: $a \neq 0$

shows $a * x^2 + b * x + c = a * (x - x_1) * (x - x_2)$

<proof>

lemma *quadratic-eq-zeroes-iff*:

assumes D: $D = b^2 - 4 * a * c$

assumes x1: $x_1 = (-b + \text{sqr}t D) / (2 * a)$

assumes x2: $x_2 = (-b - \text{sqr}t D) / (2 * a)$

assumes a: $a \neq 0$

shows $a * x^2 + b * x + c = 0 \iff (D \geq 0 \wedge (x = x_1 \vee x = x_2))$ (**is** ?z \iff -)

<proof>

lemma *quadratic-ex-zero-iff*:

$(\exists x. a * x^2 + b * x + c = 0) \iff (a \neq 0 \wedge b^2 - 4 * a * c \geq 0 \vee a = 0 \wedge (b = 0 \implies c = 0))$

for a b c::real

<proof>

lemma *Cauchy-Schwarz-eq-iff*:

shows $(\text{inner } x y)^2 = \text{inner } x x * \text{inner } y y \iff ((\exists k. x = k *_R y) \vee y = 0)$

<proof>

lemma *Cauchy-Schwarz-strict-ineq*:

$(\text{inner } x \ y)^2 < \text{inner } x \ x * \text{inner } y \ y$ **if** $y \neq 0 \wedge k. x \neq k *_R y$
 ⟨proof⟩

lemma *Cauchy-Schwarz-eq2-iff*:

$|\text{inner } x \ y| = \text{norm } x * \text{norm } y \iff ((\exists k. x = k *_R y) \vee y = 0)$
 ⟨proof⟩

lemma *Cauchy-Schwarz-strict-ineq2*:

$|\text{inner } x \ y| < \text{norm } x * \text{norm } y$ **if** $y \neq 0 \wedge k. x \neq k *_R y$
 ⟨proof⟩

lemma *gt-minus-one-absI*: $\text{abs } k < 1 \implies -1 < k$ **for** $k::\text{real}$

⟨proof⟩

lemma *gt-one-absI*: $\text{abs } k < 1 \implies k < 1$ **for** $k::\text{real}$

⟨proof⟩

lemma *abs-impossible*:

$|y1| < x1 \implies |y2| < x2 \implies x1 * x2 + y1 * y2 \neq 0$ **for** $x1 \ x2::\text{real}$
 ⟨proof⟩

lemma *vangle-eq-arctan-minus*: — TODO: generalize?!

assumes $ij: i \in \text{Basis } j \in \text{Basis}$ **and** $ij\text{-neg}: i \neq j$

assumes $xy1: |y1| < x1$

assumes $xy2: |y2| < x2$

assumes $\text{less}: y2 / x2 > y1 / x1$

shows $\text{vangle } (x1 *_R i + y1 *_R j) (x2 *_R i + y2 *_R j) = \text{arctan } (y2 / x2) - \text{arctan } (y1 / x1)$

(**is** $\text{vangle } ?u \ ?v = -$)

⟨proof⟩

lemma *vangle-le-pi2*: $0 \leq u \cdot v \implies \text{vangle } u \ v \leq \text{pi}/2$

⟨proof⟩

lemma *inner-eq-vangle*: $u \cdot v = \cos (\text{vangle } u \ v) * (\text{norm } u * \text{norm } v)$

⟨proof⟩

lemma *vangle-scaleR-self*:

$\text{vangle } (k *_R v) \ v = (\text{if } k = 0 \vee v = 0 \text{ then } \text{pi} / 2 \text{ else if } k > 0 \text{ then } 0 \text{ else } \text{pi})$

$\text{vangle } v \ (k *_R v) = (\text{if } k = 0 \vee v = 0 \text{ then } \text{pi} / 2 \text{ else if } k > 0 \text{ then } 0 \text{ else } \text{pi})$

⟨proof⟩

lemma *vangle-scaleR*:

$\text{vangle } (k *_R v) \ w = \text{vangle } v \ w \ \text{vangle } w \ (k *_R v) = \text{vangle } w \ v$ **if** $k > 0$

⟨proof⟩

lemma *cos-vangle-eq-zero-iff-vangle*:

$\cos (\text{vangle } u \ v) = 0 \iff (u = 0 \vee v = 0 \vee u \cdot v = 0)$

⟨proof⟩

lemma *ortho-imp-angle-pi-half*: $u \cdot v = 0 \implies \text{vangle } u \ v = \text{pi} / 2$
 ⟨proof⟩

lemma *arccos-eq-zero-iff*: $\text{arccos } x = 0 \iff x = 1$ **if** $-1 \leq x \leq 1$
 ⟨proof⟩

lemma *vangle-eq-zeroD*: $\text{vangle } u \ v = 0 \implies (\exists k. v = k *_{\mathbb{R}} u)$
 ⟨proof⟩

lemma *less-one-multI*:— TODO: also in AA!

fixes $e \ x :: \text{real}$

shows $e \leq 1 \implies 0 < x \implies x < 1 \implies e * x < 1$

⟨proof⟩

lemma *conemem-expansion-estimate*:

fixes $u \ v \ u' \ v' :: 'a :: \text{euclidean-space}$

assumes $t \in \{0 .. \text{pi} / 2\}$

assumes *angle-pos*: $0 < \text{vangle } u \ v \ \text{vangle } u' \ v' < \text{pi} / 2$

assumes *angle-le*: $(\text{vangle } u' \ v') \leq (\text{vangle } u \ v)$

assumes $\text{norm } u = 1 \ \text{norm } v = 1$

shows $\text{norm } (\text{conemem } u' \ v' \ t) \geq \min (\text{norm } u') (\text{norm } v') * \text{norm } (\text{conemem } u \ v \ t)$

⟨proof⟩

lemma *conemem-commute*: $\text{conemem } a \ b \ t = \text{conemem } b \ a \ (\text{pi} / 2 - t)$ **if** $0 \leq t \leq \text{pi} / 2$

⟨proof⟩

lemma *conesegment-commute*: $\text{conesegment } a \ b = \text{conesegment } b \ a$
 ⟨proof⟩

definition *conefield* $u \ v = \text{cone hull } (\text{conesegment } u \ v)$

lemma *conefield-alt-def*: $\text{conefield } u \ v = \text{cone hull } \{u - -v\}$
 ⟨proof⟩

lemma

bounded-linear-image-cone-hull:

assumes *bounded-linear* F

shows $F \text{ ` } (\text{cone hull } T) = \text{cone hull } (F \text{ ` } T)$

⟨proof⟩

lemma

bounded-linear-image-conefield:

assumes *bounded-linear* F

shows $F \text{ ` } \text{conefield } u \ v = \text{conefield } (F \ u) \ (F \ v)$

⟨proof⟩

lemma *conefield-commute*: $\text{conefield } x \ y = \text{conefield } y \ x$
 ⟨proof⟩

lemma *convex-conefield*: $\text{convex } (\text{conefield } x \ y)$
 ⟨proof⟩

lemma *conefield-scaleRI*: $v \in \text{conefield } (r *_{\mathbb{R}} x) \ y$ **if** $v \in \text{conefield } x \ y \ r > 0$
 ⟨proof⟩

lemma *conefield-scaleRD*: $v \in \text{conefield } x \ y$ **if** $v \in \text{conefield } (r *_{\mathbb{R}} x) \ y \ r > 0$
 ⟨proof⟩

lemma *conefield-scaleR*: $\text{conefield } (r *_{\mathbb{R}} x) \ y = \text{conefield } x \ y$ **if** $r > 0$
 ⟨proof⟩

lemma *conefield-expansion-estimate*:
fixes $u \ v :: 'a :: \text{euclidean-space}$ **and** $F :: 'a \Rightarrow 'a$
assumes $t \in \{0 .. \pi / 2\}$
assumes *angle-pos*: $0 < \text{vangle } u \ v < \pi / 2$
assumes *angle-le*: $\text{vangle } (F \ u) \ (F \ v) \leq \text{vangle } u \ v$
assumes *bounded-linear* F
assumes $x \in \text{conefield } u \ v$
shows $\text{norm } (F \ x) \geq \min (\text{norm } (F \ u) / \text{norm } u) (\text{norm } (F \ v) / \text{norm } v) * \text{norm } x$
 ⟨proof⟩

lemma *conefield-rightI*:
assumes *ij*: $i \in \text{Basis } j \in \text{Basis}$ **and** *ij-neq*: $i \neq j$
assumes $y \in \{y1 .. y2\}$
shows $(i + y *_{\mathbb{R}} j) \in \text{conefield } (i + y1 *_{\mathbb{R}} j) \ (i + y2 *_{\mathbb{R}} j)$
 ⟨proof⟩

lemma *conefield-right-vangleI*:
assumes *ij*: $i \in \text{Basis } j \in \text{Basis}$ **and** *ij-neq*: $i \neq j$
assumes $y \in \{y1 .. y2\} \ y1 < y2$
shows $(i + y *_{\mathbb{R}} j) \in \text{conefield } (i + y1 *_{\mathbb{R}} j) \ (i + y2 *_{\mathbb{R}} j)$
 ⟨proof⟩

lemma *cone-conefield[intro, simp]*: $\text{cone } (\text{conefield } x \ y)$
 ⟨proof⟩

lemma *conefield-mk-rightI*:
assumes *ij*: $i \in \text{Basis } j \in \text{Basis}$ **and** *ij-neq*: $i \neq j$
assumes $(i + (y / x) *_{\mathbb{R}} j) \in \text{conefield } (i + (y1 / x1) *_{\mathbb{R}} j) \ (i + (y2 / x2) *_{\mathbb{R}} j)$
assumes $x > 0 \ x1 > 0 \ x2 > 0$
shows $(x *_{\mathbb{R}} i + y *_{\mathbb{R}} j) \in \text{conefield } (x1 *_{\mathbb{R}} i + y1 *_{\mathbb{R}} j) \ (x2 *_{\mathbb{R}} i + y2 *_{\mathbb{R}} j)$
 ⟨proof⟩

```

lemma conefield-prod3I:
  assumes  $x > 0$   $x1 > 0$   $x2 > 0$ 
  assumes  $y1 / x1 \leq y / x$   $y / x \leq y2 / x2$ 
  shows  $(x, y, 0) \in (\text{conefield } (x1, y1, 0) (x2, y2, 0))::(\text{real*real*real}) \text{ set}$ 
  <proof>

end

```

7 Linear ODE

```

theory Linear-ODE

```

```

imports

```

```

  ../IVP/Flow

```

```

  Bounded-Linear-Operator

```

```

  Multivariate-Taylor

```

```

begin

```

```

lemma

```

```

  exp-scaleR-has-derivative-right[derivative-intros]:

```

```

  fixes  $f::\text{real} \Rightarrow \text{real}$ 

```

```

  assumes (f has-derivative f') (at x within s)

```

```

  shows  $((\lambda x. \text{exp } (f x *_R A)) \text{ has-derivative } (\lambda h. f' h *_R (\text{exp } (f x *_R A) * A)))$ 
  (at x within s)
  <proof>

```

```

context

```

```

fixes  $A::'a::\{\text{banach,perfect-space}\}$  blinop

```

```

begin

```

```

definition linode-solution  $t0$   $x0 = (\lambda t. \text{exp } ((t - t0) *_R A) x0)$ 

```

```

lemma linode-solution-solves-ode:

```

```

  (linode-solution  $t0$   $x0$  solves-ode  $(\lambda-. A)$ ) UNIV UNIV linode-solution  $t0$   $x0$   $t0 =$ 
   $x0$ 
  <proof>

```

```

lemma (linode-solution  $t0$   $x0$  usolves-ode  $(\lambda-. A)$  from  $t0$ ) UNIV UNIV

```

```

  <proof>

```

```

end

```

```

end

```

```

theory ODE-Analysis

```

```

imports

```

```

  Library/MVT-Ex

```

```

  IVP/Flow

```

```

  IVP/Upper-Lower-Solution

```

```

  IVP/Reachability-Analysis

```

```

  IVP/Flow-Congs

```


IVP/Cones
Library/Linear-ODE
begin

end

References

- [1] W. Walter. *Ordinary Differential Equations*. Springer, 1 edition, 1998.