

Native Words

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Abstract

This entry makes machine words and machine arithmetic available for code generation from Isabelle/HOL. It provides a common abstraction that hides the differences between the different target languages. The code generator maps these operations to the APIs of the target languages. Apart from that, we extend the available bit operations on types `int` and `integer`, and map them to the operations in the target languages.

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Chapter 1

Common logic auxiliary for all fixed-width word types

```
theory Uint-Common
imports
  HOL-Library.Word
  Word-Lib.Signed-Division-Word
  Word-Lib.Most-significant-bit
  Word-Lib.Bit-Comprehension
begin
```

1.1 Some abstract nonsense

```
lemmas [transfer-rule] =
  identity-quotient
  fun-quotient
  Quotient-integer[folded integer.pcr-cr-eq]
```

```
lemma undefined-transfer:
  assumes Quotient R Abs Rep T
  shows T (Rep undefined) undefined
  <proof>
```

```
bundle undefined-transfer = undefined-transfer[transfer-rule]
```

1.2 Establishing type class instances for type copies of word type

The lifting machinery is not localized, hence the abstract proofs are carried out using morphisms.

```
locale word-type-copy =
  fixes of-word :: <'b::len word ⇒ 'a>
  and word-of :: <'a ⇒ 'b word>
```

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```

assumes type-definition: ⟨type-definition word-of of-word UNIV⟩
begin

lemma word-of-word:
  ⟨word-of (of-word w) = w⟩
  ⟨proof⟩

lemma of-word-of [code abstype]:
  ⟨of-word (word-of p) = p⟩
  — Use an abstract type for code generation to disable pattern matching on
    of-word.
  ⟨proof⟩

lemma word-of-eqI:
  ⟨p = q⟩ if ⟨word-of p = word-of q⟩
  ⟨proof⟩

lemma eq-iff-word-of:
  ⟨p = q ⟷ word-of p = word-of q⟩
  ⟨proof⟩

end

bundle constraintless
begin

  ⟨ML⟩

end

locale word-type-copy-ring = word-type-copy
opening constraintless +
constrains word-of :: ⟨a ⇒ b::len word⟩
assumes word-of-0 [code]: ⟨word-of 0 = 0⟩
and word-of-1 [code]: ⟨word-of 1 = 1⟩
and word-of-add [code]: ⟨word-of (p + q) = word-of p + word-of q⟩
and word-of-minus [code]: ⟨word-of (- p) = - (word-of p)⟩
and word-of-diff [code]: ⟨word-of (p - q) = word-of p - word-of q⟩
and word-of-mult [code]: ⟨word-of (p * q) = word-of p * word-of q⟩
and word-of-div [code]: ⟨word-of (p div q) = word-of p div word-of q⟩
and word-of-mod [code]: ⟨word-of (p mod q) = word-of p mod word-of q⟩
and equal-iff-word-of [code]: ⟨HOL.equal p q ⟷ HOL.equal (word-of p)
  (word-of q)⟩
and less-eq-iff-word-of [code]: ⟨p ≤ q ⟷ word-of p ≤ word-of q⟩
and less-iff-word-of [code]: ⟨p < q ⟷ word-of p < word-of q⟩
begin

lemma of-class-comm-ring-1:
  ⟨OFCLASS('a, comm-ring-1-class)⟩

```

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<proof>

lemma *of-class-semiring-modulo*:

<OFCLASS('a, semiring-modulo-class)>

<proof>

lemma *of-class-equal*:

<OFCLASS('a, equal-class)>

<proof>

lemma *of-class-linorder*:

<OFCLASS('a, linorder-class)>

<proof>

end

locale *word-type-copy-bits = word-type-copy-ring*

opening *constraintless and bit-operations-syntax +*

constrains *word-of :: '<a::{comm-ring-1, semiring-modulo, equal, linorder}> =>*
'b::len word>

fixes *signed-drop-bit :: <nat => 'a => 'a>*

assumes *bit-eq-word-of [code]: <bit p = bit (word-of p)>*

and *word-of-not [code]: <word-of (NOT p) = NOT (word-of p)>*

and *word-of-and [code]: <word-of (p AND q) = word-of p AND word-of q>*

and *word-of-or [code]: <word-of (p OR q) = word-of p OR word-of q>*

and *word-of-xor [code]: <word-of (p XOR q) = word-of p XOR word-of q>*

and *word-of-mask [code]: <word-of (mask n) = mask n>*

and *word-of-push-bit [code]: <word-of (push-bit n p) = push-bit n (word-of p)>*

and *word-of-drop-bit [code]: <word-of (drop-bit n p) = drop-bit n (word-of p)>*

and *word-of-signed-drop-bit [code]: <word-of (signed-drop-bit n p) = Word.signed-drop-bit*
n (word-of p)>

and *word-of-take-bit [code]: <word-of (take-bit n p) = take-bit n (word-of p)>*

and *word-of-set-bit [code]: <word-of (Bit-Operations.set-bit n p) = Bit-Operations.set-bit*
n (word-of p)>

and *word-of-unset-bit [code]: <word-of (unset-bit n p) = unset-bit n (word-of*
p)>

and *word-of-flip-bit [code]: <word-of (flip-bit n p) = flip-bit n (word-of p)>*

begin

lemma *word-of-bool*:

<word-of (of-bool n) = of-bool n>

<proof>

lemma *word-of-nat*:

<word-of (of-nat n) = of-nat n>

<proof>

lemma *word-of-numeral [simp]*:

<word-of (numeral n) = numeral n>

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<proof>

lemma *word-of-power*:

<word-of (p ^ n) = word-of p ^ n>

<proof>

lemma *even-iff-word-of*:

<2 dvd p $\iff 2 dvd word-of p$ (is $?P \iff ?Q$)>

<proof>

lemma *of-class-ring-bit-operations*:

<OFCLASS('a, ring-bit-operations-class)>

<proof>

lemma [code]:

<take-bit n a = a AND mask n> for a :: 'a

<proof>

lemma [code]:

<mask (Suc n) = push-bit n (1 :: 'a) OR mask n>

<mask 0 = (0 :: 'a)>

<proof>

lemma [code]:

<Bit-Operations.set-bit n w = w OR push-bit n 1> for w :: 'a

<proof>

lemma [code]:

<unset-bit n w = w AND NOT (push-bit n 1)> for w :: 'a

<proof>

lemma [code]:

<flip-bit n w = w XOR push-bit n 1> for w :: 'a

<proof>

end

locale *word-type-copy-more = word-type-copy-bits +*

constrains *word-of :: 'a::{ring-bit-operations, equal, linorder} \Rightarrow 'b::len word>*

fixes *of-nat :: 'nat \Rightarrow 'a> and nat-of :: 'a \Rightarrow nat>*

and *of-int :: 'int \Rightarrow 'a> and int-of :: 'a \Rightarrow int>*

and *of-integer :: 'integer \Rightarrow 'a> and integer-of :: 'a \Rightarrow integer>*

assumes *word-of-nat-eq: <word-of (of-nat n) = word-of-nat n>*

and *nat-of-eq-word-of: <nat-of p = unat (word-of p)>*

and *word-of-int-eq: <word-of (of-int k) = word-of-int k>*

and *int-of-eq-word-of: <int-of p = uint (word-of p)>*

and *word-of-integer-eq: <word-of (of-integer l) = word-of-int (int-of-integer l)>*

and *integer-of-eq-word-of: <integer-of p = unsigned (word-of p)>*

begin

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lemma *of-word-numeral* [*code-post*]:
 $\langle \text{of-word } (\text{numeral } n) = \text{numeral } n \rangle$
 $\langle \text{proof} \rangle$

lemma *of-word-0* [*code-post*]:
 $\langle \text{of-word } 0 = 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *of-word-1* [*code-post*]:
 $\langle \text{of-word } 1 = 1 \rangle$
 $\langle \text{proof} \rangle$

Use pretty numerals from integer for pretty printing

lemma *numeral-eq-integer* [*code-unfold*]:
 $\langle \text{numeral } n = \text{of-integer } (\text{numeral } n) \rangle$
 $\langle \text{proof} \rangle$

lemma *neg-numeral-eq-integer* [*code-unfold*]:
 $\langle - \text{numeral } n = \text{of-integer } (- \text{numeral } n) \rangle$
 $\langle \text{proof} \rangle$

end

locale *word-type-copy-misc* = *word-type-copy-more*
opening *constraintless* **and** *bit-operations-syntax* +
constrains *word-of* :: $\langle 'a :: \{\text{ring-bit-operations, equal, linorder}\} \Rightarrow 'b :: \text{len word} \rangle$
fixes *size* :: *nat* **and** *set-bits-aux* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{nat} \Rightarrow 'a \Rightarrow 'a \rangle$
assumes *size-eq-length*: $\langle \text{size} = \text{LENGTH}('b :: \text{len}) \rangle$
and *msb-iff-word-of* [*code*]: $\langle \text{msb } p \longleftrightarrow \text{msb } (\text{word-of } p) \rangle$
and *size-eq-word-of*: $\langle \text{Nat.size } (p :: 'a) = \text{Nat.size } (\text{word-of } p) \rangle$
and *word-of-set-bits*: $\langle \text{word-of } (\text{set-bits } P) = \text{set-bits } P \rangle$
and *word-of-set-bits-aux*: $\langle \text{word-of } (\text{set-bits-aux } P \ n \ p) = \text{Bit-Comprehension.set-bits-aux } P \ n \ (\text{word-of } p) \rangle$
begin

lemma *size-eq* [*code*]:
 $\langle \text{Nat.size } p = \text{size} \rangle$ **for** $p :: 'a$
 $\langle \text{proof} \rangle$

lemma *set-bits-aux-code* [*code*]:
 $\langle \text{set-bits-aux } f \ n \ w =$
 $(\text{if } n = 0 \text{ then } w$
 $\text{else let } n' = n - 1 \text{ in set-bits-aux } f \ n' \ (\text{push-bit } 1 \ w \ \text{OR } (\text{if } f \ n' \ \text{then } 1 \ \text{else } 0))) \rangle$
 $\langle \text{proof} \rangle$

lemma *set-bits-code* [*code*]: $\langle \text{set-bits } P = \text{set-bits-aux } P \ \text{size } 0 \rangle$
 $\langle \text{proof} \rangle$

lemma *of-class-bit-comprehension*:
 ‹OFCLASS('a, bit-comprehension-class)›
 ‹proof›

end

1.3 Establishing operation variants tailored towards target languages

locale *word-type-copy-target-language* = *word-type-copy-misc* +
constrains *word-of* :: ‹'a::{\ring-bit-operations, equal, linorder} ⇒ 'b::len word›
fixes *size-integer* :: integer
and *almost-size* :: nat
assumes *size-integer-eq-length*: ‹size-integer = Nat.of-nat LENGTH('b::len)›
and *almost-size-eq-decr-length*: ‹almost-size = LENGTH('b::len) - Suc 0›
begin

definition *shiffl* :: ‹'a ⇒ integer ⇒ 'a›
where ‹shiffl w k = (if k < 0 ∨ size-integer ≤ k then undefined (push-bit :: nat ⇒ 'a ⇒ 'a) w k
 else push-bit (nat-of-integer k) w)›

lemma *word-of-shiffl* [code abstract]:
 ‹word-of (shiffl w k) =
 (if k < 0 ∨ size-integer ≤ k then word-of (undefined (push-bit :: - ⇒ - ⇒ 'a) w k)
 else push-bit (nat-of-integer k) (word-of w))›
 ‹proof›

lemma *push-bit-code* [code]:
 ‹push-bit k w = (if k < size then shiffl w (integer-of-nat k) else 0)›
 ‹proof›

definition *shiftr* :: ‹'a ⇒ integer ⇒ 'a›
where ‹shiftr w k = (if k < 0 ∨ size-integer ≤ k then undefined (drop-bit :: nat ⇒ 'a ⇒ 'a) w k
 else drop-bit (nat-of-integer k) w)›

lemma *word-of-shiftr* [code abstract]:
 ‹word-of (shiftr w k) =
 (if k < 0 ∨ size-integer ≤ k then word-of (undefined (drop-bit :: - ⇒ - ⇒ 'a) w k)
 else drop-bit (nat-of-integer k) (word-of w))›
 ‹proof›

lemma *drop-bit-code* [code]:
 ‹drop-bit k w = (if k < size then shiftr w (integer-of-nat k) else 0)›
 ‹proof›

definition *sshiftr* :: $\langle 'a \Rightarrow \text{integer} \Rightarrow 'a \rangle$
where $\langle \text{sshiftr } w \ k = (\text{if } k < 0 \vee \text{size-integer} \leq k \text{ then undefined (signed-drop-bit} \\ \text{:: } - \Rightarrow - \Rightarrow 'a) \ w \ k$
 $\text{else signed-drop-bit (nat-of-integer } k) \ w) \rangle$

lemma *word-of-sshiftr* [*code abstract*]:
 $\langle \text{word-of (sshiftr } w \ k) =$
 $(\text{if } k < 0 \vee \text{size-integer} \leq k \text{ then word-of (undefined (signed-drop-bit :: } - \Rightarrow - \Rightarrow$
 $'a) \ w \ k)$
 $\text{else Word.signed-drop-bit (nat-of-integer } k) \ (\text{word-of } w)) \rangle$
 $\langle \text{proof} \rangle$

lemma *signed-drop-bit-code* [*code*]:
 $\langle \text{signed-drop-bit } k \ w = (\text{if } k < \text{size} \text{ then sshiftr } w \ (\text{integer-of-nat } k)$
 $\text{else if (bit } w \ \text{almost-size) then } - \ 1 \ \text{else } 0) \rangle$
 $\langle \text{proof} \rangle$

definition *test-bit* :: $\langle 'a \Rightarrow \text{integer} \Rightarrow \text{bool} \rangle$
where $\langle \text{test-bit } w \ k = (\text{if } k < 0 \vee \text{size-integer} \leq k \text{ then undefined (bit :: } 'a \Rightarrow -)$
 $w \ k$
 $\text{else bit } w \ (\text{nat-of-integer } k)) \rangle$

lemma *test-bit-eq* [*code*]:
 $\langle \text{test-bit } w \ k = (\text{if } k < 0 \vee \text{size-integer} \leq k \text{ then undefined (bit :: } 'a \Rightarrow -) \ w \ k$
 $\text{else bit (word-of } w) \ (\text{nat-of-integer } k)) \rangle$
 $\langle \text{proof} \rangle$

lemma *bit-code* [*code*]:
 $\langle \text{bit } w \ n \iff n < \text{size} \wedge \text{test-bit } w \ (\text{integer-of-nat } n) \rangle$
 $\langle \text{proof} \rangle$

end

1.3.1 Division by signed division

Division on *'a word* is unsigned, but Scala and OCaml only have signed division and modulus.

context
begin

private lemma *div-half-nat*:
fixes $m \ n :: \text{nat}$
assumes $n \neq 0$
shows $(m \ \text{div} \ n, \ m \ \text{mod} \ n) = ($
 let
 $q = 2 * (\text{drop-bit } 1 \ m \ \text{div} \ n);$
 $r = m - q * n$
 $\text{in if } n \leq r \text{ then } (q + 1, r - n) \ \text{else } (q, r))$

```

⟨proof⟩ lemma div-half-word:
  fixes x y :: 'a :: len word
  assumes y ≠ 0
  shows (x div y, x mod y) = (
    let
      q = push-bit 1 (drop-bit 1 x div y);
      r = x - q * y
    in if y ≤ r then (q + 1, r - y) else (q, r))
⟨proof⟩

```

This algorithm implements unsigned division in terms of signed division. Taken from Hacker's Delight.

```

lemma divmod-via-sdivmod:
  fixes x y :: 'a :: len word
  assumes y ≠ 0
  shows (x div y, x mod y) = (
    if push-bit (LENGTH('a) - 1) 1 ≤ y then
      if x < y then (0, x) else (1, x - y)
    else let
      q = (push-bit 1 (drop-bit 1 x sdiv y));
      r = x - q * y
    in if y ≤ r then (q + 1, r - y) else (q, r))
⟨proof⟩

```

end

1.3.2 Conversion from *int* to *'a* word

```

lemma word-of-int-via-signed:
  includes bit-operations-syntax
  assumes shift-def: shift = push-bit LENGTH('a) 1
  and overflow-def: overflow = push-bit (LENGTH('a) - 1) 1
  shows
    (word-of-int i :: 'a :: len word) =
      (let i' = i AND mask LENGTH('a)
       in if bit i' (LENGTH('a) - 1) then
          if i' - shift < - overflow ∨ overflow ≤ i' - shift then arbitrary1 i' else
            word-of-int (i' - shift)
          else if i' < - overflow ∨ overflow ≤ i' then arbitrary2 i' else word-of-int i')
⟨proof⟩

```

1.3.3 Quickcheck conversion functions

context

includes state-combinator-syntax

begin

definition qc-random-cnv ::

(natural ⇒ 'a::term-of) ⇒ natural ⇒ Random.seed

```

    ⇒ ('a × (unit ⇒ Code-Evaluation.term)) × Random.seed
  where qc-random-cnv a-of-natural i = Random.range (i + 1) ○→ (λk. Pair (
    let n = a-of-natural k
    in (n, λ-. Code-Evaluation.term-of n)))

```

end

definition *qc-exhaustive-cnv* :: (natural ⇒ 'a) ⇒ ('a ⇒ (bool × term list) option)
 ⇒ natural ⇒ (bool × term list) option

where

```

  qc-exhaustive-cnv a-of-natural f d =
    Quickcheck-Exhaustive.exhaustive (%x. f (a-of-natural x)) d

```

definition *qc-full-exhaustive-cnv* ::

```

  (natural ⇒ ('a::term-of)) ⇒ ('a × (unit ⇒ term) ⇒ (bool × term list) option)
  ⇒ natural ⇒ (bool × term list) option

```

where

```

  qc-full-exhaustive-cnv a-of-natural f d = Quickcheck-Exhaustive.full-exhaustive
  (%(x, xt). f (a-of-natural x, %- Code-Evaluation.term-of (a-of-natural x))) d

```

declare [[quickcheck-narrowing-ghc-options = -XTypeSynonymInstances]]

definition *qc-narrowing-drawn-from* :: 'a list ⇒ integer ⇒ -

where

```

  qc-narrowing-drawn-from xs =
    foldr Quickcheck-Narrowing.sum (map Quickcheck-Narrowing.cons (butlast xs))
    (Quickcheck-Narrowing.cons (last xs))

```

locale *quickcheck-narrowing-samples* =

fixes *a-of-integer* :: integer ⇒ 'a × 'a :: {partial-term-of, term-of}

and *zero* :: 'a

and *tr* :: typerep

begin

function *narrowing-samples* :: integer ⇒ 'a list

where

```

  narrowing-samples i =
    (if i > 0 then let (a, a') = a-of-integer i in narrowing-samples (i - 1) @ [a, a']
    else [zero])
  ⟨proof⟩

```

termination including *integer.lifting*

⟨proof⟩

definition *partial-term-of-sample* :: integer ⇒ 'a

where

```

  partial-term-of-sample i =
    (if i < 0 then undefined
     else if i = 0 then zero
     else if i mod 2 = 0 then snd (a-of-integer (i div 2))

```

else fst (a-of-integer (i div 2 + 1)))

lemma *partial-term-of-code:*

partial-term-of (ty :: 'a itself) (Quickcheck-Narrowing.Narrowing-variable p t) ≡
Code-Evaluation.Free (STR "'-'') tr
partial-term-of (ty :: 'a itself) (Quickcheck-Narrowing.Narrowing-constructor i
[]) ≡
Code-Evaluation.term-of (partial-term-of-sample i)
<proof>

end

lemmas [*code*] =

quickcheck-narrowing-samples.narrowing-samples.simps
quickcheck-narrowing-samples.partial-term-of-sample-def

1.3.4 Code generator setup

code-identifier code-module *Uint-Common* \rightarrow

(SML) Word **and** *(Haskell) Word* **and** *(OCaml) Word* **and** *(Scala) Word*

end

Chapter 2

Common base for target language implementations of word types

```
theory Code-Target-Word
  imports HOL-Library.Word
begin
```

2.1 SML

The separate code target *SML-word* collects setups for the code generator that PolyML does not provide.

<ML>

2.2 Haskell

In the *Data.Bits.Bits* type class, shifts and bit indices are given as *Int* rather than *Integer*.

Additional constants take only parameters of type *integer* rather than *nat* and check the preconditions as far as possible (e.g., being non-negative) in a portable way.

```
code-printing code-module Data-Bits  $\rightarrow$  (Haskell)
```

```
<
```

```
module Data-Bits where {
```

```
import qualified Data.Bits;
```

```
{-
```

```
The ...Bounded functions assume that the Integer argument for the shift  
or bit index fits into an Int, is non-negative and (for types of fixed bit width)
```

```

    less than bitSize
  -}

infixl 7 .&;
infixl 6 'xor';
infixl 5 .|.;

(&.) :: Data.Bits.Bits a => a -> a -> a;
(&.) = (Data.Bits.&.);

xor :: Data.Bits.Bits a => a -> a -> a;
xor = Data.Bits.xor;

(|.) :: Data.Bits.Bits a => a -> a -> a;
(|.) = (Data.Bits.|.);

complement :: Data.Bits.Bits a => a -> a;
complement = Data.Bits.complement;

testBitUnbounded :: Data.Bits.Bits a => a -> Integer -> Bool;
testBitUnbounded x b
  | b <= toInteger (Prelude.maxBound :: Int) = Data.Bits.testBit x (fromInteger
b)
  | otherwise = error (Bit index too large: ++ show b)
;

testBitBounded :: Data.Bits.Bits a => a -> Integer -> Bool;
testBitBounded x b = Data.Bits.testBit x (fromInteger b);

shifflUnbounded :: Data.Bits.Bits a => a -> Integer -> a;
shifflUnbounded x n
  | n <= toInteger (Prelude.maxBound :: Int) = Data.Bits.shiftL x (fromInteger
n)
  | otherwise = error (Shift operand too large: ++ show n)
;

shifflBounded :: Data.Bits.Bits a => a -> Integer -> a;
shifflBounded x n = Data.Bits.shiftL x (fromInteger n);

shiftrUnbounded :: Data.Bits.Bits a => a -> Integer -> a;
shiftrUnbounded x n
  | n <= toInteger (Prelude.maxBound :: Int) = Data.Bits.shiftR x (fromInteger
n)
  | otherwise = error (Shift operand too large: ++ show n)
;

shiftrBounded :: (Ord a, Data.Bits.Bits a) => a -> Integer -> a;
shiftrBounded x n = Data.Bits.shiftR x (fromInteger n);

```

```
}>
```

and — *HOL.Quickcheck-Narrowing* maps *integer* to Haskell’s *Prelude.Int* type instead of *Integer*. For compatibility with the Haskell target, we nevertheless provide bounded and unbounded functions.

(*Haskell-Quickcheck*)

```
<
```

```
module Data-Bits where {
```

```
import qualified Data.Bits;
```

```
{-
```

The functions assume that the Int argument for the shift or bit index is non-negative and (for types of fixed bit width) less than bitSize

```
-}
```

```
infixl 7 .&;
```

```
infixl 6 'xor';
```

```
infixl 5 .|;
```

```
(.&.) :: Data.Bits.Bits a => a -> a -> a;
```

```
(.&.) = (Data.Bits.&.);
```

```
xor :: Data.Bits.Bits a => a -> a -> a;
```

```
xor = Data.Bits.xor;
```

```
(.|.) :: Data.Bits.Bits a => a -> a -> a;
```

```
(.|.) = (Data.Bits.|.);
```

```
complement :: Data.Bits.Bits a => a -> a;
```

```
complement = Data.Bits.complement;
```

```
testBitUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> Bool;
```

```
testBitUnbounded = Data.Bits.testBit;
```

```
testBitBounded :: Data.Bits.Bits a => a -> Prelude.Int -> Bool;
```

```
testBitBounded = Data.Bits.testBit;
```

```
shiftlUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> a;
```

```
shiftlUnbounded = Data.Bits.shiftL;
```

```
shiftlBounded :: Data.Bits.Bits a => a -> Prelude.Int -> a;
```

```
shiftlBounded = Data.Bits.shiftL;
```

```
shiftrUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> a;
```

```
shiftrUnbounded = Data.Bits.shiftR;
```

```
shiftrBounded :: (Ord a, Data.Bits.Bits a) => a -> Prelude.Int -> a;
```

```
shiftrBounded = Data.Bits.shiftR;
```

}>

code-reserved (*Haskell*) *Data-Bits*

end

Chapter 3

A special case of a conversion.

```
theory Code-Int-Integer-Conversion
imports
  Main
begin
```

Use this function to convert numeral *integers* quickly into *ints*. By default, it works only for symbolic evaluation; normally generated code raises an exception at run-time. If theory *Code-Target-Int-Bit* is imported, it works again, because then *int* is implemented in terms of *integer* even for symbolic evaluation.

```
definition int-of-integer-symbolic :: integer  $\Rightarrow$  int
  where int-of-integer-symbolic = int-of-integer
```

```
lemma int-of-integer-symbolic-aux-code [code nbe]:
  int-of-integer-symbolic 0 = 0
  int-of-integer-symbolic (Code-Numeral.Pos n) = Int.Pos n
  int-of-integer-symbolic (Code-Numeral.Neg n) = Int.Neg n
  <proof>
```

```
end
```


Chapter 4

Unsigned words of 64 bits

```
theory Uint64
imports
  HOL-Library.Code-Target-Bit-Shifts
  Uint-Common
  Code-Target-Word
  Code-Int-Integer-Conversion
begin
```

PolyML (in version 5.7) provides a `Word64` structure only when run in 64-bit mode. Therefore, we by default provide an implementation of 64-bit words using `IntInf.int` and masking. The code target `SML_word` replaces this implementation and maps the operations directly to the `Word64` structure provided by the Standard ML implementations.

The `Eval` target used by `value` and `eval` dynamically tests at runtime for the version of PolyML and uses PolyML's `Word64` structure if it detects a 64-bit version which does not suffer from a division bug found in PolyML 5.6.

4.1 Type definition and primitive operations

```
typedef uint64 =  $\langle UNIV :: 64 \text{ word set} \rangle$   $\langle proof \rangle$ 

global-interpretation uint64: word-type-copy Abs-uint64 Rep-uint64
   $\langle proof \rangle$ 

setup-lifting type-definition-uint64

declare uint64.of-word-of [code abstype]

declare Quotient-uint64 [transfer-rule]

instantiation uint64 ::  $\langle \{ comm\text{-}ring\text{-}1, semiring\text{-}modulo, equal, linorder \} \rangle$ 
begin
```

```

lift-definition zero-uint64 :: uint64 is 0 <proof>
lift-definition one-uint64 :: uint64 is 1 <proof>
lift-definition plus-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <(+)> <proof>
lift-definition uminus-uint64 :: <uint64 ⇒ uint64> is uminus <proof>
lift-definition minus-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <(-)> <proof>
lift-definition times-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <(*)> <proof>
lift-definition divide-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <(div)> <proof>
lift-definition modulo-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <(mod)> <proof>
lift-definition equal-uint64 :: <uint64 ⇒ uint64 ⇒ bool> is <HOL.equal> <proof>
lift-definition less-eq-uint64 :: <uint64 ⇒ uint64 ⇒ bool> is <(≤)> <proof>
lift-definition less-uint64 :: <uint64 ⇒ uint64 ⇒ bool> is <(<)> <proof>

```

```

global-interpretation uint64: word-type-copy-ring Abs-uint64 Rep-uint64
  <proof>

```

```

instance <proof>

```

```

end

```

```

instantiation uint64 :: ring-bit-operations
begin

```

```

lift-definition bit-uint64 :: <uint64 ⇒ nat ⇒ bool> is bit <proof>
lift-definition not-uint64 :: <uint64 ⇒ uint64> is <Bit-Operations.not> <proof>
lift-definition and-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <Bit-Operations.and>
  <proof>
lift-definition or-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <Bit-Operations.or>
  <proof>
lift-definition xor-uint64 :: <uint64 ⇒ uint64 ⇒ uint64> is <Bit-Operations.xor>
  <proof>
lift-definition mask-uint64 :: <nat ⇒ uint64> is mask <proof>
lift-definition push-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is push-bit <proof>
lift-definition drop-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is drop-bit <proof>
lift-definition signed-drop-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is signed-drop-bit
  <proof>
lift-definition take-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is take-bit <proof>
lift-definition set-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is Bit-Operations.set-bit
  <proof>
lift-definition unset-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is unset-bit <proof>
lift-definition flip-bit-uint64 :: <nat ⇒ uint64 ⇒ uint64> is flip-bit <proof>

```

```

global-interpretation uint64: word-type-copy-bits Abs-uint64 Rep-uint64 signed-drop-bit-uint64
  <proof>

```

```

instance
  <proof>

```

```

end

```

lift-definition *uint64-of-nat* :: $\langle \text{nat} \Rightarrow \text{uint64} \rangle$
is *word-of-nat* $\langle \text{proof} \rangle$

lift-definition *nat-of-uint64* :: $\langle \text{uint64} \Rightarrow \text{nat} \rangle$
is *unat* $\langle \text{proof} \rangle$

lift-definition *uint64-of-int* :: $\langle \text{int} \Rightarrow \text{uint64} \rangle$
is *word-of-int* $\langle \text{proof} \rangle$

lift-definition *int-of-uint64* :: $\langle \text{uint64} \Rightarrow \text{int} \rangle$
is *uint* $\langle \text{proof} \rangle$

context
includes *integer.lifting*
begin

lift-definition *Uint64* :: $\langle \text{integer} \Rightarrow \text{uint64} \rangle$
is *word-of-int* $\langle \text{proof} \rangle$

lift-definition *integer-of-uint64* :: $\langle \text{uint64} \Rightarrow \text{integer} \rangle$
is *uint* $\langle \text{proof} \rangle$

end

global-interpretation *uint64*: *word-type-copy-more Abs-uint64 Rep-uint64 signed-drop-bit-uint64*
uint64-of-nat nat-of-uint64 uint64-of-int int-of-uint64 Uint64 integer-of-uint64
 $\langle \text{proof} \rangle$

instantiation *uint64* :: $\{ \text{size}, \text{msb}, \text{bit-comprehension} \}$
begin

lift-definition *size-uint64* :: $\langle \text{uint64} \Rightarrow \text{nat} \rangle$ **is** *size* $\langle \text{proof} \rangle$

lift-definition *msb-uint64* :: $\langle \text{uint64} \Rightarrow \text{bool} \rangle$ **is** *msb* $\langle \text{proof} \rangle$

lift-definition *set-bits-uint64* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{uint64} \rangle$ **is** *set-bits* $\langle \text{proof} \rangle$

lift-definition *set-bits-aux-uint64* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{nat} \Rightarrow \text{uint64} \Rightarrow \text{uint64} \rangle$ **is**
set-bits-aux $\langle \text{proof} \rangle$

global-interpretation *uint64*: *word-type-copy-misc Abs-uint64 Rep-uint64 signed-drop-bit-uint64*
uint64-of-nat nat-of-uint64 uint64-of-int int-of-uint64 Uint64 integer-of-uint64 64
set-bits-aux-uint64
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

4.2 Code setup

For SML, we generate an implementation of unsigned 64-bit words using `IntInf.int`. If `LargeWord.wordSize > 63` of the Isabelle/ML runtime environment holds, then we assume that there is also a `Word64` structure available and accordingly replace the implementation for the target `Eval`.

code-printing code-module `Uint64` \rightarrow (SML) \langle (* Test that words can handle numbers between 0 and 63 *)

`val - = if 64 <= Word.wordSize then () else raise (Fail (wordSize less than 6));`

```

structure Uint64 : sig
  eqtype uint64;
  val zero : uint64;
  val one : uint64;
  val fromInt : IntInf.int -> uint64;
  val toInt : uint64 -> IntInf.int;
  val toLarge : uint64 -> LargeWord.word;
  val fromLarge : LargeWord.word -> uint64
  val plus : uint64 -> uint64 -> uint64;
  val minus : uint64 -> uint64 -> uint64;
  val times : uint64 -> uint64 -> uint64;
  val divide : uint64 -> uint64 -> uint64;
  val modulus : uint64 -> uint64 -> uint64;
  val negate : uint64 -> uint64;
  val less-eq : uint64 -> uint64 -> bool;
  val less : uint64 -> uint64 -> bool;
  val notb : uint64 -> uint64;
  val andb : uint64 -> uint64 -> uint64;
  val orb : uint64 -> uint64 -> uint64;
  val xorb : uint64 -> uint64 -> uint64;
  val shiffl : uint64 -> IntInf.int -> uint64;
  val shiftr : uint64 -> IntInf.int -> uint64;
  val shiftr-signed : uint64 -> IntInf.int -> uint64;
  val test-bit : uint64 -> IntInf.int -> bool;
end = struct

type uint64 = IntInf.int;

val mask = 0xFFFFFFFFFFFFFFFF : IntInf.int;

val zero = 0 : IntInf.int;

val one = 1 : IntInf.int;

fun fromInt x = IntInf.andb(x, mask);

fun toInt x = x

```

```

fun toLarge x = LargeWord.fromLargeInt (IntInf.toLarge x);
fun fromLarge x = IntInf.fromLarge (LargeWord.toLargeInt x);
fun plus x y = IntInf.andb(IntInf.+(x, y), mask);
fun minus x y = IntInf.andb(IntInf.-(x, y), mask);
fun negate x = IntInf.andb(IntInf.~(x), mask);
fun times x y = IntInf.andb(IntInf.*(x, y), mask);
fun divide x y = IntInf.div(x, y);
fun modulus x y = IntInf.mod(x, y);
fun less-eq x y = IntInf.<=(x, y);
fun less x y = IntInf.<(x, y);
fun notb x = IntInf.andb(IntInf.notb(x), mask);
fun orb x y = IntInf.orb(x, y);
fun andb x y = IntInf.andb(x, y);
fun xorb x y = IntInf.xorb(x, y);
val maxWord = IntInf.pow (2, Word.wordSize);

fun shiftl x n =
  if n < maxWord then IntInf.andb(IntInf.<<(x, Word.fromLargeInt (IntInf.toLarge
n)), mask)
  else 0;

fun shiftr x n =
  if n < maxWord then IntInf.~>>(x, Word.fromLargeInt (IntInf.toLarge n))
  else 0;

val msb-mask = 0x8000000000000000 : IntInf.int;

fun shiftr-signed x i =
  if IntInf.andb(x, msb-mask) = 0 then shiftr x i
  else if i >= 64 then 0xFFFFFFFFFFFFFFFF
  else let
    val x' = shiftr x i
    val m' = IntInf.andb(IntInf.<<(mask, Word.max(0w64 - Word.fromLargeInt
(IntInf.toLarge i), 0w0)), mask)
  in IntInf.orb(x', m') end;

```

```

fun test-bit x n =
  if n < maxWord then IntInf.andb (x, IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge
n))) <> 0
  else false;

end
>
code-reserved (SML) Uint64

```

<ML>

```

code-printing code-module Uint64 → (Haskell)
<module Uint64 (Int64, Word64) where

```

```

  import Data.Int(Int64)
  import Data.Word(Word64)
code-reserved (Haskell) Uint64

```

OCaml and Scala provide only signed 64bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

```

code-printing code-module Uint64 → (OCaml)

```

```

<module Uint64 : sig
  val less : int64 -> int64 -> bool
  val less-eq : int64 -> int64 -> bool
  val shiffl : int64 -> Z.t -> int64
  val shiftr : int64 -> Z.t -> int64
  val shiftr-signed : int64 -> Z.t -> int64
  val test-bit : int64 -> Z.t -> bool
end = struct

```

```

(* negative numbers have their highest bit set,
   so they are greater than positive ones *)

```

```

let less x y =
  if Int64.compare x Int64.zero < 0 then
    Int64.compare y Int64.zero < 0 && Int64.compare x y < 0
  else Int64.compare y Int64.zero < 0 || Int64.compare x y < 0;;

```

```

let less-eq x y =
  if Int64.compare x Int64.zero < 0 then
    Int64.compare y Int64.zero < 0 && Int64.compare x y <= 0
  else Int64.compare y Int64.zero < 0 || Int64.compare x y <= 0;;

```

```

let shiffl x n = Int64.shift-left x (Z.to-int n);;

```

```

let shiftr x n = Int64.shift-right-logical x (Z.to-int n);;

```

```

let shiftr-signed x n = Int64.shift-right x (Z.to-int n);;

```

```

let test-bit x n =
  Int64.compare
    (Int64.logand x (Int64.shift-left Int64.one (Z.to-int n)))
    Int64.zero
  <> 0;;

end;; (*struct Uint64*)
code-reserved (OCaml) Uint64

code-printing code-module Uint64 → (Scala)
object Uint64 {

def less(x: Long, y: Long) : Boolean =
  x < 0 match {
    case true => y < 0 && x < y
    case false => y < 0 || x < y
  }

def less-eq(x: Long, y: Long) : Boolean =
  x < 0 match {
    case true => y < 0 && x <= y
    case false => y < 0 || x <= y
  }

def shiftl(x: Long, n: BigInt) : Long = x << n.intValue

def shiftr(x: Long, n: BigInt) : Long = x >>> n.intValue

def shiftr-signed(x: Long, n: BigInt) : Long = x >> n.intValue

def test-bit(x: Long, n: BigInt) : Boolean =
  (x & (1L << n.intValue)) != 0

} /* object Uint64 */
code-reserved (Scala) Uint64

```

OCaml's conversion from `Big_int` to `int64` demands that the value fits into a signed 64-bit integer. The following justifies the implementation.

context

includes *bit-operations-syntax*

begin

definition *Uint64-signed* :: integer ⇒ uint64

where *Uint64-signed* $i =$ (if $i < -(0x8000000000000000) \vee i \geq 0x8000000000000000$ then undefined *Uint64* i else *Uint64* i)

lemma *Uint64-code* [code]:

Uint64 $i =$
 (let $i' = i$ AND $0xFFFFFFFFFFFFFFFF$

in if bit i' 63 then *Uint64-signed* ($i' - 0x1000000000000000$) else *Uint64-signed* i')

including *undefined-transfer* **and** *integer.lifting* \langle proof \rangle

lemma *Uint64-signed-code* [code]:

Rep-uint64 (*Uint64-signed* i) =
 (if $i < -(0x8000000000000000) \vee i \geq 0x8000000000000000$ then *Rep-uint64*
 (*undefined Uint64* i) else *word-of-int* (*int-of-integer-symbolic* i))
 \langle proof \rangle

end

Avoid *Abs-uint64* in generated code, use *Rep-uint64'* instead. The symbolic implementations for *code_simp* use *Rep-uint64*.

The new destructor *Rep-uint64'* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint64* directly, because that makes *code_simp* loop.

If code generation raises *Match*, some equation probably contains *Rep-uint64* ([code abstract] equations for *uint64* may use *Rep-uint64* because these instances will be folded away.)

To convert *64 word* values into *uint64*, use *Abs-uint64'*.

definition *Rep-uint64'* **where** [simp]: *Rep-uint64'* = *Rep-uint64*

lemma *Rep-uint64'-transfer* [transfer-rule]:

rel-fun cr-uint64 (=) ($\lambda x. x$) *Rep-uint64'*
 \langle proof \rangle

lemma *Rep-uint64'-code* [code]: *Rep-uint64'* $x = (BITS\ n.\ bit\ x\ n)$

\langle proof \rangle

lift-definition *Abs-uint64'* :: *64 word* \Rightarrow *uint64* **is** $\lambda x :: 64\ word.\ x$ \langle proof \rangle

lemma *Abs-uint64'-code* [code]:

Abs-uint64' $x = Uint64$ (*integer-of-int* (*uint* x))

including *integer.lifting* \langle proof \rangle

declare [[code drop: *term-of-class.term-of* :: *uint64* \Rightarrow -]]

lemma *term-of-uint64-code* [code]:

defines *TR* \equiv *typerep.Typeprep* **and** *bit0* \equiv *STR "Numeral-Type.bit0"*

shows

term-of-class.term-of $x =$

Code-Evaluation.App (*Code-Evaluation.Const* (*STR "Uint64.uint64.Abs-uint64"*)
 (*TR* (*STR "fun"*) [*TR* (*STR "Word.word"*) [*TR bit0* [*TR bit0* [*TR bit0* [*TR bit0*
 [*TR bit0* [*TR bit0* [*TR* (*STR "Numeral-Type.num1"*) []]]]]]]], *TR* (*STR "Uint64.uint64"*)
 []))

(*term-of-class.term-of* (*Rep-uint64'* x))

\langle proof \rangle

code-printing

```

type-constructor uint64  $\rightarrow$ 
  (SML) Uint64.uint64 and
  (Haskell) Uint64.Word64 and
  (OCaml) int64 and
  (Scala) Long
| constant Uint64  $\rightarrow$ 
  (SML) Uint64.fromInt and
  (Haskell) (Prelude.fromInteger - :: Uint64.Word64) and
  (Haskell-Quickcheck) (Prelude.fromInteger (Prelude.toInteger -) :: Uint64.Word64)
and
  (Scala) -.longValue
| constant Uint64-signed  $\rightarrow$ 
  (OCaml) Z.to'-int64
| constant 0 :: uint64  $\rightarrow$ 
  (SML) Uint64.zero and
  (Haskell) (0 :: Uint64.Word64) and
  (OCaml) Int64.zero and
  (Scala) 0
| constant 1 :: uint64  $\rightarrow$ 
  (SML) Uint64.one and
  (Haskell) (1 :: Uint64.Word64) and
  (OCaml) Int64.one and
  (Scala) 1
| constant plus :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.plus and
  (Haskell) infixl 6 + and
  (OCaml) Int64.add and
  (Scala) infixl 7 +
| constant uminus :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.negate and
  (Haskell) negate and
  (OCaml) Int64.neg and
  (Scala) !(-)
| constant minus :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.minus and
  (Haskell) infixl 6 - and
  (OCaml) Int64.sub and
  (Scala) infixl 7 -
| constant times :: uint64  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.times and
  (Haskell) infixl 7 * and
  (OCaml) Int64.mul and
  (Scala) infixl 8 *
| constant HOL.equal :: uint64  $\Rightarrow$  -  $\Rightarrow$  bool  $\rightarrow$ 
  (SML) !((- : Uint64.uint64) = -) and
  (Haskell) infix 4 == and
  (OCaml) (Int64.compare - - = 0) and

```

```

(Scala) infixl 5 ==
| class-instance uint64 :: equal  $\rightarrow$ 
  (Haskell) -
| constant less-eq :: uint64  $\Rightarrow$  -  $\Rightarrow$  bool  $\rightarrow$ 
  (SML) Uint64.less'-eq and
  (Haskell) infix 4 <= and
  (OCaml) Uint64.less'-eq and
  (Scala) Uint64.less'-eq
| constant less :: uint64  $\Rightarrow$  -  $\Rightarrow$  bool  $\rightarrow$ 
  (SML) Uint64.less and
  (Haskell) infix 4 < and
  (OCaml) Uint64.less and
  (Scala) Uint64.less
| constant Bit-Operations.not :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.notb and
  (Haskell) Data'-Bits.complement and
  (OCaml) Int64.lognot and
  (Scala) -.unary'~
| constant Bit-Operations.and :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.andb and
  (Haskell) infixl 7 Data-Bits.&. and
  (OCaml) Int64.logand and
  (Scala) infixl 3 &
| constant Bit-Operations.or :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.orb and
  (Haskell) infixl 5 Data-Bits.|. and
  (OCaml) Int64.logor and
  (Scala) infixl 1 |
| constant Bit-Operations.xor :: uint64  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Uint64.xorb and
  (Haskell) Data'-Bits.xor and
  (OCaml) Int64.logxor and
  (Scala) infixl 2 ^

definition uint64-divmod :: uint64  $\Rightarrow$  uint64  $\Rightarrow$  uint64  $\times$  uint64 where
  uint64-divmod x y =
    (if y = 0 then (undefined ((div) :: uint64  $\Rightarrow$  -) x (0 :: uint64), undefined ((mod)
:: uint64  $\Rightarrow$  -) x (0 :: uint64)))
    else (x div y, x mod y))

definition uint64-div :: uint64  $\Rightarrow$  uint64  $\Rightarrow$  uint64
where uint64-div x y = fst (uint64-divmod x y)

definition uint64-mod :: uint64  $\Rightarrow$  uint64  $\Rightarrow$  uint64
where uint64-mod x y = snd (uint64-divmod x y)

lemma div-uint64-code [code]: x div y = (if y = 0 then 0 else uint64-div x y)
including undefined-transfer <proof>

```

lemma *mod-uint64-code* [code]: $x \text{ mod } y = (\text{if } y = 0 \text{ then } x \text{ else } \text{uint64-mod } x \ y)$
including *undefined-transfer* ⟨proof⟩

definition *uint64-sdiv* :: $\text{uint64} \Rightarrow \text{uint64} \Rightarrow \text{uint64}$
where [code del]:
uint64-sdiv $x \ y =$
 ($\text{if } y = 0 \text{ then undefined } ((\text{div}) :: \text{uint64} \Rightarrow -) \ x \ (0 :: \text{uint64})$
 else $\text{Abs-uint64 } (\text{Rep-uint64 } \ x \ \text{sdiv } \text{Rep-uint64 } \ y)$)

definition *div0-uint64* :: $\text{uint64} \Rightarrow \text{uint64}$
where [code del]: $\text{div0-uint64 } \ x = \text{undefined } ((\text{div}) :: \text{uint64} \Rightarrow -) \ x \ (0 :: \text{uint64})$
declare [[code abort: *div0-uint64*]]

definition *mod0-uint64* :: $\text{uint64} \Rightarrow \text{uint64}$
where [code del]: $\text{mod0-uint64 } \ x = \text{undefined } ((\text{mod}) :: \text{uint64} \Rightarrow -) \ x \ (0 :: \text{uint64})$
declare [[code abort: *mod0-uint64*]]

lemma *uint64-divmod-code* [code]:
uint64-divmod $x \ y =$
 ($\text{if } 0x8000000000000000 \leq y \text{ then if } x < y \text{ then } (0, \ x) \text{ else } (1, \ x - y)$
 else $\text{if } y = 0 \text{ then } (\text{div0-uint64 } \ x, \ \text{mod0-uint64 } \ x)$
 else $\text{let } q = \text{push-bit } 1 \ (\text{uint64-sdiv } (\text{drop-bit } 1 \ x) \ y);$
 $r = x - q * y$
 $\text{in if } r \geq y \text{ then } (q + 1, \ r - y) \text{ else } (q, \ r)$)
including *undefined-transfer* ⟨proof⟩

lemma *uint64-sdiv-code* [code]:
 $\text{Rep-uint64 } (\text{uint64-sdiv } \ x \ y) =$
 ($\text{if } y = 0 \text{ then } \text{Rep-uint64 } (\text{undefined } ((\text{div}) :: \text{uint64} \Rightarrow -) \ x \ (0 :: \text{uint64}))$
 else $\text{Rep-uint64 } \ x \ \text{sdiv } \text{Rep-uint64 } \ y$)
 ⟨proof⟩

Note that we only need a translation for signed division, but not for the remainder because $\text{uint64-divmod } \ ?x \ ?y = (\text{if } 9223372036854775808 \leq ?y \text{ then if } ?x < ?y \text{ then } (0, \ ?x) \text{ else } (1, \ ?x - ?y) \text{ else if } ?y = 0 \text{ then } (\text{div0-uint64 } \ ?x, \ \text{mod0-uint64 } \ ?x) \text{ else let } q = \text{push-bit } 1 \ (\text{uint64-sdiv } (\text{drop-bit } 1 \ ?x) \ ?y); \ r = ?x - q * ?y \text{ in if } ?y \leq r \text{ then } (q + 1, \ r - ?y) \text{ else } (q, \ r))$ computes both with division only.

code-printing

constant *uint64-div* \rightarrow
 (SML) *Uint64.divide* **and**
 (Haskell) *Prelude.div*
| **constant** *uint64-mod* \rightarrow
 (SML) *Uint64.modulus* **and**
 (Haskell) *Prelude.mod*
| **constant** *uint64-divmod* \rightarrow
 (Haskell) *divmod*
| **constant** *uint64-sdiv* \rightarrow
 (OCaml) *Int64.div* **and**

(Scala) - '/ -

global-interpretation *uint64*: *word-type-copy-target-language Abs-uint64 Rep-uint64 signed-drop-bit-uint64 uint64-of-nat nat-of-uint64 uint64-of-int int-of-uint64 Uint64 integer-of-uint64 64 set-bits-aux-uint64 64 63*

defines *uint64-test-bit* = *uint64.test-bit*
and *uint64-shiffl* = *uint64.shiffl*
and *uint64-shiftr* = *uint64.shiftr*
and *uint64-sshiftr* = *uint64.sshiftr*
 ⟨*proof*⟩

code-printing constant *uint64-test-bit* \rightarrow

(SML) *Uint64.test'-bit* **and**
 (Haskell) *Data'-Bits.testBitBounded* **and**
 (OCaml) *Uint64.test'-bit* **and**
 (Scala) *Uint64.test'-bit* **and**
 (Eval) (fn x => fn i => if i < 0 orelse i >= 64 then raise (Fail argument to *uint64'-test'-bit* out of bounds) else *Uint64.test'-bit* x i)

code-printing constant *uint64-shiffl* \rightarrow

(SML) *Uint64.shiffl* **and**
 (Haskell) *Data'-Bits.shifflBounded* **and**
 (OCaml) *Uint64.shiffl* **and**
 (Scala) *Uint64.shiffl* **and**
 (Eval) (fn w => fn i => if i < 0 orelse i >= 64 then raise (Fail argument to *uint64'-shiffl* out of bounds) else *Uint64.shiffl* w i)

code-printing constant *uint64-shiftr* \rightarrow

(SML) *Uint64.shiftr* **and**
 (Haskell) *Data'-Bits.shiftrBounded* **and**
 (OCaml) *Uint64.shiftr* **and**
 (Scala) *Uint64.shiftr* **and**
 (Eval) (fn w => fn i => if i < 0 orelse i >= 64 then raise (Fail argument to *uint64'-shiftr* out of bounds) else *Uint64.shiftr* w i)

code-printing constant *uint64-sshiftr* \rightarrow

(SML) *Uint64.shiftr'-signed* **and**
 (Haskell)
 (*Prelude.fromInteger (Prelude.toInteger (Data'-Bits.shiftrBounded (Prelude.fromInteger (Prelude.toInteger -) :: Uint64.Int64) -)) :: Uint64.Word64*) **and**
 (OCaml) *Uint64.shiftr'-signed* **and**
 (Scala) *Uint64.shiftr'-signed* **and**
 (Eval) (fn w => fn i => if i < 0 orelse i >= 64 then raise (Fail argument to *uint64'-shiftr'-signed* out of bounds) else *Uint64.shiftr'-signed* w i)

context

includes *bit-operations-syntax*
begin

lemma *uint64-msb-test-bit*: $msb\ x \longleftrightarrow bit\ (x :: uint64)\ 63$
 ⟨proof⟩

lemma *msb-uint64-code* [code]: $msb\ x \longleftrightarrow uint64\text{-test-bit}\ x\ 63$
 ⟨proof⟩

lemma *uint64-of-int-code* [code]:
 $uint64\text{-of-int}\ i = Uint64\ (integer\text{-of-int}\ i)$
including *integer.lifting* ⟨proof⟩

lemma *int-of-uint64-code* [code]:
 $int\text{-of-uint64}\ x = int\text{-of-integer}\ (integer\text{-of-uint64}\ x)$
including *integer.lifting* ⟨proof⟩

lemma *uint64-of-nat-code* [code]:
 $uint64\text{-of-nat} = uint64\text{-of-int} \circ int$
 ⟨proof⟩

lemma *nat-of-uint64-code* [code]:
 $nat\text{-of-uint64}\ x = nat\text{-of-integer}\ (integer\text{-of-uint64}\ x)$
 ⟨proof⟩ **including** *integer.lifting* ⟨proof⟩

definition *integer-of-uint64-signed* :: $uint64 \Rightarrow integer$
where

$integer\text{-of-uint64}\text{-signed}\ n = (if\ bit\ n\ 63\ then\ undefined\ integer\text{-of-uint64}\ n\ else\ integer\text{-of-uint64}\ n)$

lemma *integer-of-uint64-signed-code* [code]:
 $integer\text{-of-uint64}\text{-signed}\ n =$
 (if bit n 63 then undefined integer-of-uint64 n else integer-of-int (uint (Rep-uint64' n)))
 ⟨proof⟩

lemma *integer-of-uint64-code* [code]:
 $integer\text{-of-uint64}\ n =$
 (if bit n 63 then integer-of-uint64-signed (n AND 0x7FFFFFFFFFFFFFFF) OR 0x8000000000000000 else integer-of-uint64-signed n)
 ⟨proof⟩
including *integer.lifting* ⟨proof⟩

end

code-printing

constant *integer-of-uint64* \mapsto
 (SML) *Uint64.toInt* **and**
 (Haskell) *Prelude.toInteger*
 | **constant** *integer-of-uint64-signed* \mapsto
 (OCaml) *Z.of'-int64* **and**

(Scala) *BigInt*

4.3 Quickcheck setup

```

definition wint64-of-natural :: natural ⇒ wint64
where wint64-of-natural x ≡ Uint64 (integer-of-natural x)

instantiation wint64 :: {random, exhaustive, full-exhaustive} begin
definition random-wint64 ≡ qc-random-cnv wint64-of-natural
definition exhaustive-wint64 ≡ qc-exhaustive-cnv wint64-of-natural
definition full-exhaustive-wint64 ≡ qc-full-exhaustive-cnv wint64-of-natural
instance ⟨proof⟩
end

instantiation wint64 :: narrowing begin

interpretation quickcheck-narrowing-samples
  λi. let x = Uint64 i in (x, 0xFFFFFFFFFFFFFFFF - x) 0
  Typerep.Typerep (STR "Uint64.wint64") [] ⟨proof⟩

definition narrowing-wint64 d = qc-narrowing-drawn-from (narrowing-samples d)
d
declare [[code drop: partial-term-of :: wint64 itself ⇒ -]]
lemmas partial-term-of-wint64 [code] = partial-term-of-code

instance ⟨proof⟩
end

end

```

Chapter 5

Unsigned words of 32 bits

```
theory Uint32
  imports
    HOL-Library.Code-Target-Bit-Shifts
    Uint-Common
    Code-Target-Word
    Code-Int-Integer-Conversion
begin

5.1 Type definition and primitive operations

typedef uint32 = ⟨UNIV :: 32 word set⟩ ⟨proof⟩

global-interpretation uint32: word-type-copy Abs-uint32 Rep-uint32
  ⟨proof⟩

setup-lifting type-definition-uint32

declare uint32.of-word-of [code abstype]

declare Quotient-uint32 [transfer-rule]

instantiation uint32 :: ⟨{comm-ring-1, semiring-modulo, equal, linorder}⟩
begin

lift-definition zero-uint32 :: uint32 is 0 ⟨proof⟩
lift-definition one-uint32 :: uint32 is 1 ⟨proof⟩
lift-definition plus-uint32 :: ⟨uint32 ⇒ uint32 ⇒ uint32⟩ is ⟨(+)⟩ ⟨proof⟩
lift-definition uminus-uint32 :: ⟨uint32 ⇒ uint32⟩ is uminus ⟨proof⟩
lift-definition minus-uint32 :: ⟨uint32 ⇒ uint32 ⇒ uint32⟩ is ⟨(-)⟩ ⟨proof⟩
lift-definition times-uint32 :: ⟨uint32 ⇒ uint32 ⇒ uint32⟩ is ⟨(*)⟩ ⟨proof⟩
lift-definition divide-uint32 :: ⟨uint32 ⇒ uint32 ⇒ uint32⟩ is ⟨(div)⟩ ⟨proof⟩
lift-definition modulo-uint32 :: ⟨uint32 ⇒ uint32 ⇒ uint32⟩ is ⟨(mod)⟩ ⟨proof⟩
lift-definition equal-uint32 :: ⟨uint32 ⇒ uint32 ⇒ bool⟩ is ⟨HOL.equal⟩ ⟨proof⟩
lift-definition less-eq-uint32 :: ⟨uint32 ⇒ uint32 ⇒ bool⟩ is ⟨(≤)⟩ ⟨proof⟩
```

lift-definition *less-uint32* :: $\langle \text{uint32} \Rightarrow \text{uint32} \Rightarrow \text{bool} \rangle$ **is** $\langle (<) \rangle$ $\langle \text{proof} \rangle$

global-interpretation *uint32*: *word-type-copy-ring Abs-uint32 Rep-uint32*
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

instantiation *uint32* :: *ring-bit-operations*
begin

lift-definition *bit-uint32* :: $\langle \text{uint32} \Rightarrow \text{nat} \Rightarrow \text{bool} \rangle$ **is** *bit* $\langle \text{proof} \rangle$

lift-definition *not-uint32* :: $\langle \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** $\langle \text{Bit-Operations.not} \rangle$ $\langle \text{proof} \rangle$

lift-definition *and-uint32* :: $\langle \text{uint32} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** $\langle \text{Bit-Operations.and} \rangle$
 $\langle \text{proof} \rangle$

lift-definition *or-uint32* :: $\langle \text{uint32} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** $\langle \text{Bit-Operations.or} \rangle$
 $\langle \text{proof} \rangle$

lift-definition *xor-uint32* :: $\langle \text{uint32} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** $\langle \text{Bit-Operations.xor} \rangle$
 $\langle \text{proof} \rangle$

lift-definition *mask-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \rangle$ **is** *mask* $\langle \text{proof} \rangle$

lift-definition *push-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** *push-bit* $\langle \text{proof} \rangle$

lift-definition *drop-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** *drop-bit* $\langle \text{proof} \rangle$

lift-definition *signed-drop-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** *signed-drop-bit*
 $\langle \text{proof} \rangle$

lift-definition *take-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** *take-bit* $\langle \text{proof} \rangle$

lift-definition *set-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** $\langle \text{Bit-Operations.set-bit} \rangle$
 $\langle \text{proof} \rangle$

lift-definition *unset-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** *unset-bit* $\langle \text{proof} \rangle$

lift-definition *flip-bit-uint32* :: $\langle \text{nat} \Rightarrow \text{uint32} \Rightarrow \text{uint32} \rangle$ **is** *flip-bit* $\langle \text{proof} \rangle$

global-interpretation *uint32*: *word-type-copy-bits Abs-uint32 Rep-uint32 signed-drop-bit-uint32*
 $\langle \text{proof} \rangle$

instance
 $\langle \text{proof} \rangle$

end

lift-definition *uint32-of-nat* :: $\langle \text{nat} \Rightarrow \text{uint32} \rangle$
is *word-of-nat* $\langle \text{proof} \rangle$

lift-definition *nat-of-uint32* :: $\langle \text{uint32} \Rightarrow \text{nat} \rangle$
is *unat* $\langle \text{proof} \rangle$

lift-definition *uint32-of-int* :: $\langle \text{int} \Rightarrow \text{uint32} \rangle$
is *word-of-int* $\langle \text{proof} \rangle$

lift-definition *int-of-uint32* :: $\langle \text{uint32} \Rightarrow \text{int} \rangle$

```

  is uint ⟨proof⟩

context
  includes integer.lifting
begin

lift-definition Uint32 :: ⟨integer ⇒ uint32⟩
  is word-of-int ⟨proof⟩

lift-definition integer-of-uint32 :: ⟨uint32 ⇒ integer⟩
  is uint ⟨proof⟩

end

global-interpretation uint32: word-type-copy-more Abs-uint32 Rep-uint32 signed-drop-bit-uint32
  uint32-of-nat nat-of-uint32 uint32-of-int int-of-uint32 Uint32 integer-of-uint32
  ⟨proof⟩

instantiation uint32 :: {size, msb, bit-comprehension}
begin

lift-definition size-uint32 :: ⟨uint32 ⇒ nat⟩ is size ⟨proof⟩

lift-definition msb-uint32 :: ⟨uint32 ⇒ bool⟩ is msb ⟨proof⟩

lift-definition set-bits-uint32 :: ⟨(nat ⇒ bool) ⇒ uint32⟩ is set-bits ⟨proof⟩
lift-definition set-bits-aux-uint32 :: ⟨(nat ⇒ bool) ⇒ nat ⇒ uint32 ⇒ uint32⟩ is
  set-bits-aux ⟨proof⟩

global-interpretation uint32: word-type-copy-misc Abs-uint32 Rep-uint32 signed-drop-bit-uint32
  uint32-of-nat nat-of-uint32 uint32-of-int int-of-uint32 Uint32 integer-of-uint32 32
  set-bits-aux-uint32
  ⟨proof⟩

instance ⟨proof⟩

end

```

5.2 Code setup

```

code-printing code-module Uint32 ↪ (SML)
⟨(* Test that words can handle numbers between 0 and 31 *)
val - = if 5 <= Word.wordSize then () else raise (Fail (wordSize less than 5));

structure Uint32 : sig
  val shiftl : Word32.word -> IntInf.int -> Word32.word
  val shiftr : Word32.word -> IntInf.int -> Word32.word
  val shiftr-signed : Word32.word -> IntInf.int -> Word32.word
  val test-bit : Word32.word -> IntInf.int -> bool

```

```

end = struct

fun shiftl x n =
  Word32.<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
  Word32.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
  Word32.~>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
  Word32.andb (x, Word32.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n)))
  <> Word32.fromInt 0

end; (* struct Uint32 *)
code-reserved (SML) Uint32

code-printing code-module Uint32 → (Haskell)
  <module Uint32(Int32, Word32) where

  import Data.Int(Int32)
  import Data.Word(Word32)
code-reserved (Haskell) Uint32

```

OCaml and Scala provide only signed 32bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

```

code-printing code-module Uint32 → (OCaml)
  <module Uint32 : sig
    val less : int32 -> int32 -> bool
    val less-eq : int32 -> int32 -> bool
    val shiftl : int32 -> Z.t -> int32
    val shiftr : int32 -> Z.t -> int32
    val shiftr-signed : int32 -> Z.t -> int32
    val test-bit : int32 -> Z.t -> bool
  end = struct

  (* negative numbers have their highest bit set,
     so they are greater than positive ones *)
  let less x y =
    if Int32.compare x Int32.zero < 0 then
      Int32.compare y Int32.zero < 0 && Int32.compare x y < 0
    else Int32.compare y Int32.zero < 0 || Int32.compare x y < 0;;

  let less-eq x y =
    if Int32.compare x Int32.zero < 0 then
      Int32.compare y Int32.zero < 0 && Int32.compare x y <= 0
    else Int32.compare y Int32.zero < 0 || Int32.compare x y <= 0;;

```

```

let shiftl x n = Int32.shift-left x (Z.to-int n);

let shiftr x n = Int32.shift-right-logical x (Z.to-int n);

let shiftr-signed x n = Int32.shift-right x (Z.to-int n);

let test-bit x n =
  Int32.compare
    (Int32.logand x (Int32.shift-left Int32.one (Z.to-int n)))
    Int32.zero
  <> 0;;

```

```

end;; (*struct Uint32*)
code-reserved (OCaml) Uint32

```

```

code-printing code-module Uint32  $\rightarrow$  (Scala)

```

```

<object Uint32 {

```

```

def less(x: Int, y: Int) : Boolean =
  x < 0 match {
    case true => y < 0 && x < y
    case false => y < 0 || x < y
  }

```

```

def less-eq(x: Int, y: Int) : Boolean =
  x < 0 match {
    case true => y < 0 && x <= y
    case false => y < 0 || x <= y
  }

```

```

def shiftl(x: Int, n: BigInt) : Int = x << n.toIntValue

```

```

def shiftr(x: Int, n: BigInt) : Int = x >>> n.toIntValue

```

```

def shiftr-signed(x: Int, n: BigInt) : Int = x >> n.toIntValue

```

```

def test-bit(x: Int, n: BigInt) : Boolean =
  (x & (1 << n.toIntValue)) != 0

```

```

} /* object Uint32 */

```

```

code-reserved (Scala) Uint32

```

OCaml's conversion from `Big_int` to `int32` demands that the value fits into a signed 32-bit integer. The following justifies the implementation.

```

context

```

```

  includes bit-operations-syntax

```

```

begin

```

```

definition Uint32-signed :: integer  $\Rightarrow$  uint32

```

where *Uint32-signed* $i = (if\ i < -(0x80000000) \vee i \geq 0x80000000\ then\ undefined\ Uint32\ i\ else\ Uint32\ i)$

lemma *Uint32-code* [code]:

Uint32 $i =$
 (let $i' = i\ AND\ 0xFFFFFFFF$
 in if bit $i'\ 31$ then *Uint32-signed* ($i' - 0x10000000$) else *Uint32-signed* i')
including *undefined-transfer* and *integer.lifting* <proof>

lemma *Uint32-signed-code* [code]:

Rep-uint32 (*Uint32-signed* i) =
 (if $i < -(0x80000000) \vee i \geq 0x80000000$ then *Rep-uint32* (*undefined* *Uint32* i)
 else *word-of-int* (*int-of-integer-symbolic* i))
 <proof>

end

Avoid *Abs-uint32* in generated code, use *Rep-uint32'* instead. The symbolic implementations for *code_simp* use *Rep-uint32*.

The new destructor *Rep-uint32'* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint32* directly, because that makes *code_simp* loop.

If code generation raises *Match*, some equation probably contains *Rep-uint32* ([code abstract] equations for *uint32* may use *Rep-uint32* because these instances will be folded away.)

To convert *32 word* values into *uint32*, use *Abs-uint32'*.

definition *Rep-uint32'* **where** [simp]: *Rep-uint32'* = *Rep-uint32*

lemma *Rep-uint32'-transfer* [transfer-rule]:

rel-fun *cr-uint32* (=) ($\lambda x. x$) *Rep-uint32'*
 <proof>

lemma *Rep-uint32'-code* [code]: *Rep-uint32'* $x = (BITS\ n.\ bit\ x\ n)$

<proof>

lift-definition *Abs-uint32'* :: *32 word* \Rightarrow *uint32* **is** $\lambda x :: 32\ word.\ x$ <proof>

lemma *Abs-uint32'-code* [code]:

Abs-uint32' $x = Uint32\ (integer-of-int\ (uint\ x))$

including *integer.lifting* <proof>

declare [[code drop: *term-of-class.term-of* :: *uint32* \Rightarrow -]]

lemma *term-of-uint32-code* [code]:

defines *TR* $\equiv typerep.Type$ rep and *bit0* $\equiv STR\ "Numeral-Type.bit0"$

shows

term-of-class.term-of $x =$

Code-Evaluation.App (*Code-Evaluation.Const* (*STR* "*Uint32.uint32.Abs-uint32'*"))

```
(TR (STR "fun") [TR (STR "Word.word") [TR bit0 [TR bit0 [TR bit0 [TR bit0
[TR bit0 [TR (STR "Numeral-Type.num1") []]]]]], TR (STR "Uint32.uint32")
[]))
  (term-of-class.term-of (Rep-uint32' x))
⟨proof⟩
```

code-printing

```
type-constructor uint32 →
  (SML) Word32.word and
  (Haskell) Uint32.Word32 and
  (OCaml) int32 and
  (Scala) Int and
  (Eval) Word32.word
| constant Uint32 →
  (SML) Word32.fromLargeInt (IntInf.toLarge -) and
  (Haskell) (Prelude.fromInteger - :: Uint32.Word32) and
  (Haskell-Quickcheck) (Prelude.fromInteger (Prelude.toInteger -) :: Uint32.Word32)
and
  (Scala) -.intValue
| constant Uint32-signed →
  (OCaml) Z.to'-int32
| constant 0 :: uint32 →
  (SML) (Word32.fromInt 0) and
  (Haskell) (0 :: Uint32.Word32) and
  (OCaml) Int32.zero and
  (Scala) 0
| constant 1 :: uint32 →
  (SML) (Word32.fromInt 1) and
  (Haskell) (1 :: Uint32.Word32) and
  (OCaml) Int32.one and
  (Scala) 1
| constant plus :: uint32 ⇒ - →
  (SML) Word32.+ ((-), (-)) and
  (Haskell) infixl 6 + and
  (OCaml) Int32.add and
  (Scala) infixl 7 +
| constant uminus :: uint32 ⇒ - →
  (SML) Word32.~ and
  (Haskell) negate and
  (OCaml) Int32.neg and
  (Scala) !(-)
| constant minus :: uint32 ⇒ - →
  (SML) Word32.- ((-), (-)) and
  (Haskell) infixl 6 - and
  (OCaml) Int32.sub and
  (Scala) infixl 7 -
| constant times :: uint32 ⇒ - ⇒ - →
  (SML) Word32.* ((-), (-)) and
  (Haskell) infixl 7 * and
```

```

(OCaml) Int32.mul and
(Scala) infixl 8 *
| constant HOL.equal :: uint32 ⇒ - ⇒ bool →
(SML) !((- : Word32.word) = -) and
(Haskell) infix 4 == and
(OCaml) (Int32.compare - - = 0) and
(Scala) infixl 5 ==
| class-instance uint32 :: equal →
(Haskell) -
| constant less-eq :: uint32 ⇒ - ⇒ bool →
(SML) Word32.<= ((-), (-)) and
(Haskell) infix 4 <= and
(OCaml) Uint32.less'-eq and
(Scala) Uint32.less'-eq
| constant less :: uint32 ⇒ - ⇒ bool →
(SML) Word32.< ((-), (-)) and
(Haskell) infix 4 < and
(OCaml) Uint32.less and
(Scala) Uint32.less
| constant Bit-Operations.not :: uint32 ⇒ - →
(SML) Word32.notb and
(Haskell) Data'-Bits.complement and
(OCaml) Int32.lognot and
(Scala) -.unary'~
| constant Bit-Operations.and :: uint32 ⇒ - →
(SML) Word32.andb ((-),/ (-)) and
(Haskell) infixl 7 Data'-Bits.&. and
(OCaml) Int32.logand and
(Scala) infixl 3 &
| constant Bit-Operations.or :: uint32 ⇒ - →
(SML) Word32.orb ((-),/ (-)) and
(Haskell) infixl 5 Data'-Bits.|. and
(OCaml) Int32.logor and
(Scala) infixl 1 |
| constant Bit-Operations.xor :: uint32 ⇒ - →
(SML) Word32.xorb ((-),/ (-)) and
(Haskell) Data'-Bits.xor and
(OCaml) Int32.logxor and
(Scala) infixl 2 ^

definition uint32-divmod :: uint32 ⇒ uint32 ⇒ uint32 × uint32 where
  uint32-divmod x y =
    (if y = 0 then (undefined ((div) :: uint32 ⇒ -) x (0 :: uint32), undefined ((mod)
:: uint32 ⇒ -) x (0 :: uint32))
    else (x div y, x mod y))

```

```

definition uint32-div :: uint32 ⇒ uint32 ⇒ uint32
where uint32-div x y = fst (uint32-divmod x y)

```

definition *uint32-mod* :: *uint32* \Rightarrow *uint32* \Rightarrow *uint32*
where *uint32-mod* *x y* = *snd* (*uint32-divmod* *x y*)

lemma *div-uint32-code* [*code*]: *x div y* = (*if y = 0 then 0 else uint32-div* *x y*)
including *undefined-transfer* *<proof>*

lemma *mod-uint32-code* [*code*]: *x mod y* = (*if y = 0 then x else uint32-mod* *x y*)
including *undefined-transfer* *<proof>*

definition *uint32-sdiv* :: *uint32* \Rightarrow *uint32* \Rightarrow *uint32*
where [*code del*]:
uint32-sdiv *x y* =
 (*if y = 0 then undefined* ((*div*) :: *uint32* \Rightarrow -) *x* (*0* :: *uint32*)
else Abs-uint32 (*Rep-uint32* *x sdiv Rep-uint32 y*)

definition *div0-uint32* :: *uint32* \Rightarrow *uint32*
where [*code del*]: *div0-uint32* *x* = *undefined* ((*div*) :: *uint32* \Rightarrow -) *x* (*0* :: *uint32*)
declare [[*code abort: div0-uint32*]]

definition *mod0-uint32* :: *uint32* \Rightarrow *uint32*
where [*code del*]: *mod0-uint32* *x* = *undefined* ((*mod*) :: *uint32* \Rightarrow -) *x* (*0* :: *uint32*)
declare [[*code abort: mod0-uint32*]]

lemma *uint32-divmod-code* [*code*]:
uint32-divmod *x y* =
 (*if 0x80000000* \leq *y* then *if x < y* then (*0*, *x*) *else* (*1*, *x - y*)
else if y = 0 then (*div0-uint32* *x*, *mod0-uint32* *x*)
else let q = push-bit 1 (*uint32-sdiv* (*drop-bit 1* *x*) *y*);
 *r = x - q * y*
in if r \geq y then (*q + 1*, *r - y*) *else* (*q*, *r*)
including *undefined-transfer* *<proof>*

lemma *uint32-sdiv-code* [*code*]:
Rep-uint32 (*uint32-sdiv* *x y*) =
 (*if y = 0* then *Rep-uint32* (*undefined* ((*div*) :: *uint32* \Rightarrow -) *x* (*0* :: *uint32*))
else Rep-uint32 *x sdiv Rep-uint32 y*)
<proof>

Note that we only need a translation for signed division, but not for the remainder because *uint32-divmod* *?x ?y* = (*if 2147483648* \leq *?y* then *if ?x < ?y* then (*0*, *?x*) *else* (*1*, *?x - ?y*) *else if ?y = 0* then (*div0-uint32* *?x*, *mod0-uint32* *?x*) *else let q = push-bit 1* (*uint32-sdiv* (*drop-bit 1* *?x*) *?y*); *r = ?x - q * ?y in if ?y \leq r* then (*q + 1*, *r - ?y*) *else* (*q*, *r*) computes both with division only.

code-printing
constant *uint32-div* \rightarrow
 (*SML*) *Word32.div* ((-), (-)) **and**
 (*Haskell*) *Prelude.div*
| **constant** *uint32-mod* \rightarrow

```

(SML) Word32.mod ((-), (-)) and
(Haskell) Prelude.mod
| constant uint32-divmod  $\rightarrow$ 
(Haskell) divmod
| constant uint32-sdiv  $\rightarrow$ 
(OCaml) Int32.div and
(Scala) - / -

```

global-interpretation *uint32*: *word-type-copy-target-language Abs-uint32 Rep-uint32 signed-drop-bit-uint32 uint32-of-nat nat-of-uint32 uint32-of-int int-of-uint32 Uint32 integer-of-uint32 32 set-bits-aux-uint32 32 31*

```

defines uint32-test-bit = uint32.test-bit
  and uint32-shiffl = uint32.shiffl
  and uint32-shiftr = uint32.shiftr
  and uint32-sshiftr = uint32.sshiftr
⟨proof⟩

```

code-printing constant *uint32-test-bit* \rightarrow

```

(SML) Uint32.test'-bit and
(Haskell) Data'-Bits.testBitBounded and
(OCaml) Uint32.test'-bit and
(Scala) Uint32.test'-bit and
(Eval) (fn w => fn n => if n < 0 orelse 32 <= n then raise (Fail argument to
uint32'-test'-bit out of bounds) else Uint32.test'-bit w n)

```

code-printing constant *uint32-shiffl* \rightarrow

```

(SML) Uint32.shiffl and
(Haskell) Data'-Bits.shifflBounded and
(OCaml) Uint32.shiffl and
(Scala) Uint32.shiffl and
(Eval) (fn w => fn i => if i < 0 orelse i >= 32 then raise Fail argument to
uint32'-shiffl out of bounds else Uint32.shiffl w i)

```

code-printing constant *uint32-shiftr* \rightarrow

```

(SML) Uint32.shiftr and
(Haskell) Data'-Bits.shiftrBounded and
(OCaml) Uint32.shiftr and
(Scala) Uint32.shiftr and
(Eval) (fn w => fn i => if i < 0 orelse i >= 32 then raise Fail argument to
uint32'-shiftr out of bounds else Uint32.shiftr w i)

```

code-printing constant *uint32-sshiftr* \rightarrow

```

(SML) Uint32.shiftr'-signed and
(Haskell)
(Prelude.fromInteger (Prelude.toInteger (Data'-Bits.shiftrBounded (Prelude.fromInteger
(Prelude.toInteger -) :: Uint32.Int32) -)) :: Uint32.Word32) and
(OCaml) Uint32.shiftr'-signed and
(Scala) Uint32.shiftr'-signed and

```

(Eval) (fn w => fn i => if i < 0 orelse i >= 32 then raise Fail argument to uint32'-shiftr'-signed out of bounds else Uint32.shiftr'-signed w i)

context

includes *bit-operations-syntax*

begin

lemma *uint32-msb-test-bit*: $msb\ x \longleftrightarrow bit\ (x :: uint32)\ 31$
 ⟨proof⟩

lemma *msb-uint32-code* [code]: $msb\ x \longleftrightarrow uint32-test-bit\ x\ 31$
 ⟨proof⟩

lemma *uint32-of-int-code* [code]:
 $uint32-of-int\ i = Uint32\ (integer-of-int\ i)$
including *integer.lifting* ⟨proof⟩

lemma *int-of-uint32-code* [code]:
 $int-of-uint32\ x = int-of-integer\ (integer-of-uint32\ x)$
including *integer.lifting* ⟨proof⟩

lemma *uint32-of-nat-code* [code]:
 $uint32-of-nat = uint32-of-int \circ int$
 ⟨proof⟩

lemma *nat-of-uint32-code* [code]:
 $nat-of-uint32\ x = nat-of-integer\ (integer-of-uint32\ x)$
 ⟨proof⟩ **including** *integer.lifting* ⟨proof⟩

definition *integer-of-uint32-signed* :: $uint32 \Rightarrow integer$

where

$integer-of-uint32-signed\ n = (if\ bit\ n\ 31\ then\ undefined\ integer-of-uint32\ n\ else\ integer-of-uint32\ n)$

lemma *integer-of-uint32-signed-code* [code]:
 $integer-of-uint32-signed\ n =$
 (if bit n 31 then undefined integer-of-uint32 n else integer-of-int (uint (Rep-uint32' n)))
 ⟨proof⟩

lemma *integer-of-uint32-code* [code]:
 $integer-of-uint32\ n =$
 (if bit n 31 then integer-of-uint32-signed (n AND 0x7FFFFFFF) OR 0x80000000
 else integer-of-uint32-signed n)
 ⟨proof⟩
including *integer.lifting* ⟨proof⟩

end

code-printing

```

constant integer-of-uint32  $\rightarrow$ 
  (SML) IntInf.fromLarge (Word32.toLargeInt -) : IntInf.int and
  (Haskell) Prelude.toInteger
| constant integer-of-uint32-signed  $\rightarrow$ 
  (OCaml) Z.of'-int32 and
  (Scala) BigInt

```

5.3 Quickcheck setup

```

definition uint32-of-natural :: natural  $\Rightarrow$  uint32
where uint32-of-natural x  $\equiv$  Uint32 (integer-of-natural x)

```

```

instantiation uint32 :: {random, exhaustive, full-exhaustive} begin
definition random-uint32  $\equiv$  qc-random-cnv uint32-of-natural
definition exhaustive-uint32  $\equiv$  qc-exhaustive-cnv uint32-of-natural
definition full-exhaustive-uint32  $\equiv$  qc-full-exhaustive-cnv uint32-of-natural
instance  $\langle$ proof $\rangle$ 
end

```

```

instantiation uint32 :: narrowing begin

```

```

interpretation quickcheck-narrowing-samples
   $\lambda i.$  let x = Uint32 i in (x, 0xFFFFFFFF - x) 0
  Typerep.Typerep (STR "Uint32.uint32") []  $\langle$ proof $\rangle$ 

```

```

definition narrowing-uint32 d = qc-narrowing-drawn-from (narrowing-samples d)
d

```

```

declare [[code drop: partial-term-of :: uint32 itself  $\Rightarrow$  -]]
lemmas partial-term-of-uint32 [code] = partial-term-of-code

```

```

instance  $\langle$ proof $\rangle$ 
end

```

```

end

```

Chapter 6

Unsigned words of 16 bits

```
theory Uint16
imports
  HOL-Library.Code-Target-Bit-Shifts
  Uint-Common
  Code-Target-Word
  Code-Int-Integer-Conversion
begin
```

Restriction for ML code generation: This theory assumes that the ML system provides a `Word16` implementation (mlton does, but PolyML 5.5 does not). Therefore, the code setup lives in the target *SML-word* rather than *SML*. This ensures that code generation still works as long as *uint16* is not involved. For the target *SML* itself, no special code generation for this type is set up. Nevertheless, it should work by emulation via *16 word* if the theory *Code-Target-Int-Bit* is imported.

Restriction for OCaml code generation: OCaml does not provide an `int16` type, so no special code generation for this type is set up.

6.1 Type definition and primitive operations

```
typedef uint16 =  $\langle UNIV :: 16 \text{ word set} \rangle$   $\langle proof \rangle$ 

global-interpretation uint16: word-type-copy Abs-uint16 Rep-uint16
   $\langle proof \rangle$ 

setup-lifting type-definition-uint16

declare uint16.of-word-of [code abstype]

declare Quotient-uint16 [transfer-rule]

instantiation uint16 ::  $\langle \{ comm\text{-ring-1, semiring-modulo, equal, linorder} \} \rangle$ 
begin
```

```

lift-definition zero-uint16 :: uint16 is 0 <proof>
lift-definition one-uint16 :: uint16 is 1 <proof>
lift-definition plus-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <(+)> <proof>
lift-definition uminus-uint16 :: <uint16 ⇒ uint16> is uminus <proof>
lift-definition minus-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <(-)> <proof>
lift-definition times-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <(*)> <proof>
lift-definition divide-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <(div)> <proof>
lift-definition modulo-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <(mod)> <proof>
lift-definition equal-uint16 :: <uint16 ⇒ uint16 ⇒ bool> is <HOL.equal> <proof>
lift-definition less-eq-uint16 :: <uint16 ⇒ uint16 ⇒ bool> is <(≤)> <proof>
lift-definition less-uint16 :: <uint16 ⇒ uint16 ⇒ bool> is <(<)> <proof>

```

```

global-interpretation uint16: word-type-copy-ring Abs-uint16 Rep-uint16
  <proof>

```

```

instance <proof>

```

```

end

```

```

instantiation uint16 :: ring-bit-operations
begin

```

```

lift-definition bit-uint16 :: <uint16 ⇒ nat ⇒ bool> is bit <proof>
lift-definition not-uint16 :: <uint16 ⇒ uint16> is <Bit-Operations.not> <proof>
lift-definition and-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <Bit-Operations.and>
  <proof>
lift-definition or-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <Bit-Operations.or>
  <proof>
lift-definition xor-uint16 :: <uint16 ⇒ uint16 ⇒ uint16> is <Bit-Operations.xor>
  <proof>
lift-definition mask-uint16 :: <nat ⇒ uint16> is mask <proof>
lift-definition push-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is push-bit <proof>
lift-definition drop-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is drop-bit <proof>
lift-definition signed-drop-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is signed-drop-bit
  <proof>
lift-definition take-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is take-bit <proof>
lift-definition set-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is Bit-Operations.set-bit
  <proof>
lift-definition unset-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is unset-bit <proof>
lift-definition flip-bit-uint16 :: <nat ⇒ uint16 ⇒ uint16> is flip-bit <proof>

```

```

global-interpretation uint16: word-type-copy-bits Abs-uint16 Rep-uint16 signed-drop-bit-uint16
  <proof>

```

```

instance
  <proof>

```

```

end

```

lift-definition *uint16-of-nat* :: $\langle \text{nat} \Rightarrow \text{uint16} \rangle$
is *word-of-nat* $\langle \text{proof} \rangle$

lift-definition *nat-of-uint16* :: $\langle \text{uint16} \Rightarrow \text{nat} \rangle$
is *unat* $\langle \text{proof} \rangle$

lift-definition *uint16-of-int* :: $\langle \text{int} \Rightarrow \text{uint16} \rangle$
is *word-of-int* $\langle \text{proof} \rangle$

lift-definition *int-of-uint16* :: $\langle \text{uint16} \Rightarrow \text{int} \rangle$
is *uint* $\langle \text{proof} \rangle$

context
includes *integer.lifting*
begin

lift-definition *Uint16* :: $\langle \text{integer} \Rightarrow \text{uint16} \rangle$
is *word-of-int* $\langle \text{proof} \rangle$

lift-definition *integer-of-uint16* :: $\langle \text{uint16} \Rightarrow \text{integer} \rangle$
is *uint* $\langle \text{proof} \rangle$

end

global-interpretation *uint16*: *word-type-copy-more Abs-uint16 Rep-uint16 signed-drop-bit-uint16*
uint16-of-nat nat-of-uint16 uint16-of-int int-of-uint16 Uint16 integer-of-uint16
 $\langle \text{proof} \rangle$

instantiation *uint16* :: $\{ \text{size}, \text{msb}, \text{bit-comprehension} \}$
begin

lift-definition *size-uint16* :: $\langle \text{uint16} \Rightarrow \text{nat} \rangle$ **is** *size* $\langle \text{proof} \rangle$

lift-definition *msb-uint16* :: $\langle \text{uint16} \Rightarrow \text{bool} \rangle$ **is** *msb* $\langle \text{proof} \rangle$

lift-definition *set-bits-uint16* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{uint16} \rangle$ **is** *set-bits* $\langle \text{proof} \rangle$

lift-definition *set-bits-aux-uint16* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{nat} \Rightarrow \text{uint16} \Rightarrow \text{uint16} \rangle$ **is**
set-bits-aux $\langle \text{proof} \rangle$

global-interpretation *uint16*: *word-type-copy-misc Abs-uint16 Rep-uint16 signed-drop-bit-uint16*
uint16-of-nat nat-of-uint16 uint16-of-int int-of-uint16 Uint16 integer-of-uint16 16
set-bits-aux-uint16
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

6.2 Code setup

```

code-printing code-module Uint16  $\rightarrow$  (SML-word)
⟨(* Test that words can handle numbers between 0 and 15 *)
val - = if 4 <= Word.wordSize then () else raise (Fail (wordSize less than 4));

structure Uint16 : sig
  val shiftl : Word16.word  $\rightarrow$  IntInf.int  $\rightarrow$  Word16.word
  val shiftr : Word16.word  $\rightarrow$  IntInf.int  $\rightarrow$  Word16.word
  val shiftr-signed : Word16.word  $\rightarrow$  IntInf.int  $\rightarrow$  Word16.word
  val test-bit : Word16.word  $\rightarrow$  IntInf.int  $\rightarrow$  bool
end = struct

fun shiftl x n =
  Word16.<<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
  Word16.>>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
  Word16.~>>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
  Word16.andb (x, Word16.<<< (0wx1, Word.fromLargeInt (IntInf.toLarge n)))
<> Word16.fromInt 0

end; (* struct Uint16 *)
code-reserved (SML-word) Uint16

code-printing code-module Uint16  $\rightarrow$  (Haskell)
⟨module Uint16(Int16, Word16) where

  import Data.Int(Int16)
  import Data.Word(Word16)
code-reserved (Haskell) Uint16

```

Scala provides unsigned 16-bit numbers as Char.

```

code-printing code-module Uint16  $\rightarrow$  (Scala)
⟨object Uint16 {

  def shiftl(x: scala.Char, n: BigInt) : scala.Char = (x <<< n.intValue).toChar

  def shiftr(x: scala.Char, n: BigInt) : scala.Char = (x >>> n.intValue).toChar

  def shiftr-signed(x: scala.Char, n: BigInt) : scala.Char = (x.toShort >> n.intValue).toChar

  def test-bit(x: scala.Char, n: BigInt) : Boolean = (x & (1.toChar <<< n.intValue))
  != 0

```

```

} /* object Uint16 */
code-reserved (Scala) Uint16

```

Avoid *Abs-uint16* in generated code, use *Rep-uint16'* instead. The symbolic implementations for `code_simp` use *Rep-uint16*.

The new destructor *Rep-uint16'* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint16* directly, because that makes `code_simp` loop.

If code generation raises `Match`, some equation probably contains *Rep-uint16* ([code abstract] equations for *uint16* may use *Rep-uint16* because these instances will be folded away.)

To convert *16 word* values into *uint16*, use *Abs-uint16'*.

definition *Rep-uint16'* **where** [simp]: *Rep-uint16' = Rep-uint16*

lemma *Rep-uint16'-transfer* [transfer-rule]:
rel-fun cr-uint16 (=) (λx. x) Rep-uint16'
 ⟨proof⟩

lemma *Rep-uint16'-code* [code]: *Rep-uint16' x = (BITS n. bit x n)*
 ⟨proof⟩

lift-definition *Abs-uint16' :: 16 word ⇒ uint16 is λx :: 16 word. x* ⟨proof⟩

lemma *Abs-uint16'-code* [code]:
Abs-uint16' x = Uint16 (integer-of-int (uint x))
including *integer.lifting* ⟨proof⟩

declare [[code drop: term-of-class.term-of :: uint16 ⇒ -]]

lemma *term-of-uint16-code* [code]:
defines *TR* ≡ *typerep.TypeRep* **and** *bit0* ≡ *STR "Numeral-Type.bit0"* **shows**
term-of-class.term-of x =
Code-Evaluation.App (Code-Evaluation.Const (STR "Uint16.uint16.Abs-uint16"))
(TR (STR "fun") [TR (STR "Word.word") [TR bit0 [TR bit0 [TR bit0 [TR bit0
[TR (STR "Numeral-Type.num1") []]]], TR (STR "Uint16.uint16") []])
(term-of-class.term-of (Rep-uint16' x))
 ⟨proof⟩

lemma *Uint16-code* [code]: *Rep-uint16 (Uint16 i) = word-of-int (int-of-integer-symbolic i)*
 ⟨proof⟩

code-printing

```

type-constructor uint16 ↪
  (SML-word) Word16.word and
  (Haskell) Uint16.Word16 and
  (Scala) scala.Char
| constant Uint16 ↪

```

```

(SML-word) Word16.fromLargeInt (IntInf.toLarge -) and
(Haskell) (Prelude.fromInteger - :: Uint16.Word16) and
(Haskell-Quickcheck) (Prelude.fromInteger (Prelude.toInteger -) :: Uint16.Word16)
and
(Scala) -.charValue
| constant 0 :: uint16 →
(SML-word) (Word16.fromInt 0) and
(Haskell) (0 :: Uint16.Word16) and
(Scala) 0
| constant 1 :: uint16 →
(SML-word) (Word16.fromInt 1) and
(Haskell) (1 :: Uint16.Word16) and
(Scala) 1
| constant plus :: uint16 ⇒ - ⇒ - →
(SML-word) Word16.+ ((-), (-)) and
(Haskell) infixl 6 + and
(Scala) (- +/ -).toChar
| constant uminus :: uint16 ⇒ - →
(SML-word) Word16.~ and
(Haskell) negate and
(Scala) (- -).toChar
| constant minus :: uint16 ⇒ - →
(SML-word) Word16.- ((-), (-)) and
(Haskell) infixl 6 - and
(Scala) (- -/ -).toChar
| constant times :: uint16 ⇒ - ⇒ - →
(SML-word) Word16.* ((-), (-)) and
(Haskell) infixl 7 * and
(Scala) (- */ -).toChar
| constant HOL.equal :: uint16 ⇒ - ⇒ bool →
(SML-word) !((- : Word16.word) = -) and
(Haskell) infix 4 == and
(Scala) infixl 5 ==
| class-instance uint16 :: equal → (Haskell) -
| constant less-eq :: uint16 ⇒ - ⇒ bool →
(SML-word) Word16.<= ((-), (-)) and
(Haskell) infix 4 <= and
(Scala) infixl 4 <=
| constant less :: uint16 ⇒ - ⇒ bool →
(SML-word) Word16.< ((-), (-)) and
(Haskell) infix 4 < and
(Scala) infixl 4 <
| constant Bit-Operations.not :: uint16 ⇒ - →
(SML-word) Word16.notb and
(Haskell) Data'-Bits.complement and
(Scala) -.unary'~.toChar
| constant Bit-Operations.and :: uint16 ⇒ - →
(SML-word) Word16.andb ((-), (-)) and
(Haskell) infixl 7 Data-Bits.&. and

```

```

(Scala) (- & -).toChar
| constant Bit-Operations.or :: uint16 ⇒ - →
(SML-word) Word16.orb ((-),/ (-)) and
(Haskell) infixl 5 Data-Bits.|. and
(Scala) (- | -).toChar
| constant Bit-Operations.xor :: uint16 ⇒ - →
(SML-word) Word16.xorb ((-),/ (-)) and
(Haskell) Data'-Bits.xor and
(Scala) (- ^ -).toChar

```

definition *uint16-div* :: *uint16* ⇒ *uint16* ⇒ *uint16*
where *uint16-div* *x y* = (if *y* = 0 then undefined ((*div*) :: *uint16* ⇒ -) *x* (0 :: *uint16*) else *x div y*)

definition *uint16-mod* :: *uint16* ⇒ *uint16* ⇒ *uint16*
where *uint16-mod* *x y* = (if *y* = 0 then undefined ((*mod*) :: *uint16* ⇒ -) *x* (0 :: *uint16*) else *x mod y*)

context includes *undefined-transfer* **begin**

lemma *div-uint16-code* [*code*]: *x div y* = (if *y* = 0 then 0 else *uint16-div* *x y*)
⟨*proof*⟩

lemma *mod-uint16-code* [*code*]: *x mod y* = (if *y* = 0 then *x* else *uint16-mod* *x y*)
⟨*proof*⟩

lemma *uint16-div-code* [*code*]:
Rep-uint16 (*uint16-div* *x y*) =
(if *y* = 0 then *Rep-uint16* (undefined ((*div*) :: *uint16* ⇒ -) *x* (0 :: *uint16*)) else
Rep-uint16 *x div Rep-uint16 y*)
⟨*proof*⟩

lemma *uint16-mod-code* [*code*]:
Rep-uint16 (*uint16-mod* *x y*) =
(if *y* = 0 then *Rep-uint16* (undefined ((*mod*) :: *uint16* ⇒ -) *x* (0 :: *uint16*)) else
Rep-uint16 *x mod Rep-uint16 y*)
⟨*proof*⟩

end

code-printing constant *uint16-div* →
(SML-word) *Word16.div* ((-), (-)) **and**
(Haskell) *Prelude.div* **and**
(Scala) (- '/ -).toChar
| **constant** *uint16-mod* →
(SML-word) *Word16.mod* ((-), (-)) **and**
(Haskell) *Prelude.mod* **and**
(Scala) (- '% -).toChar

global-interpretation *uint16*: *word-type-copy-target-language Abs-uint16 Rep-uint16 signed-drop-bit-uint16*

uint16-of-nat nat-of-uint16 uint16-of-int int-of-uint16 Uint16 integer-of-uint16 16 set-bits-aux-uint16 16 15

defines *uint16-test-bit* = *uint16.test-bit*
and *uint16-shiffl* = *uint16.shiffl*
and *uint16-shiftr* = *uint16.shiftr*
and *uint16-sshiftr* = *uint16.sshiftr*
<proof>

code-printing constant *uint16-test-bit* \rightarrow
(SML-word) Uint16.test'-bit **and**
(Haskell) Data'-Bits.testBitBounded **and**
(Scala) Uint16.test'-bit

code-printing constant *uint16-shiffl* \rightarrow
(SML-word) Uint16.shiffl **and**
(Haskell) Data'-Bits.shifflBounded **and**
(Scala) Uint16.shiffl

code-printing constant *uint16-shiftr* \rightarrow
(SML-word) Uint16.shiftr **and**
(Haskell) Data'-Bits.shiftrBounded **and**
(Scala) Uint16.shiftr

code-printing constant *uint16-sshiftr* \rightarrow
(SML-word) Uint16.shiftr'-signed **and**
(Haskell)
(Prelude.fromInteger (Prelude.toInteger (Data'-Bits.shiftrBounded (Prelude.fromInteger (Prelude.toInteger -) :: Uint16.Int16) -)) :: Uint16.Word16) **and**
(Scala) Uint16.shiftr'-signed

lemma *uint16-msb-test-bit*: *msb x \longleftrightarrow bit (x :: uint16) 15*
<proof>

lemma *msb-uint16-code* [code]: *msb x \longleftrightarrow uint16-test-bit x 15*
<proof>

lemma *uint16-of-int-code* [code]: *uint16-of-int i = Uint16 (integer-of-int i)*
including *integer.lifting* *<proof>*

lemma *int-of-uint16-code* [code]:
int-of-uint16 x = int-of-integer (integer-of-uint16 x)
<proof>

lemma *uint16-of-nat-code* [code]:
uint16-of-nat = uint16-of-int \circ int
<proof>

lemma *nat-of-uint16-code* [code]:
nat-of-uint16 $x = \text{nat-of-integer } (\text{integer-of-uint16 } x)$
 ⟨proof⟩ **including** *integer.lifting* ⟨proof⟩

lemma *integer-of-uint16-code* [code]:
integer-of-uint16 $n = \text{integer-of-int } (\text{uint } (\text{Rep-uint16}' n))$
 ⟨proof⟩

code-printing

constant *integer-of-uint16* \rightarrow
 (SML-word) *Word16.toInt* - : *IntInf.int* **and**
 (Haskell) *Prelude.toInteger* **and**
 (Scala) *BigInt*

6.3 Quickcheck setup

definition *uint16-of-natural* :: *natural* \Rightarrow *uint16*
where *uint16-of-natural* $x \equiv \text{Uint16 } (\text{integer-of-natural } x)$

instantiation *uint16* :: {*random, exhaustive, full-exhaustive*} **begin**

definition *random-uint16* $\equiv \text{qc-random-cnv } \text{uint16-of-natural}$

definition *exhaustive-uint16* $\equiv \text{qc-exhaustive-cnv } \text{uint16-of-natural}$

definition *full-exhaustive-uint16* $\equiv \text{qc-full-exhaustive-cnv } \text{uint16-of-natural}$

instance ⟨proof⟩

end

instantiation *uint16* :: *narrowing* **begin**

interpretation *quickcheck-narrowing-samples*

$\lambda i. \text{let } x = \text{Uint16 } i \text{ in } (x, 0xFFFF - x) \quad 0$

Typerep.Typerep (STR "Uint16.uint16") [] ⟨proof⟩

definition *narrowing-uint16* $d = \text{qc-narrowing-drawn-from } (\text{narrowing-samples } d)$
 d

declare [[code drop: *partial-term-of* :: *uint16* *itself* \Rightarrow -]]

lemmas *partial-term-of-uint16* [code] = *partial-term-of-code*

instance ⟨proof⟩

end

end

Chapter 7

Unsigned words of 8 bits

```
theory Uint8
  imports
    HOL-Library.Code-Target-Bit-Shifts
    Uint-Common
    Code-Target-Word
    Code-Int-Integer-Conversion
begin
```

Restriction for OCaml code generation: OCaml does not provide an `int8` type, so no special code generation for this type is set up. If the theory `Code-Target-Int-Bit` is imported, the type `uint8` is emulated via `8 word`.

7.1 Type definition and primitive operations

```
typedef uint8 = ⟨UNIV :: 8 word set⟩ ⟨proof⟩

global-interpretation uint8: word-type-copy Abs-uint8 Rep-uint8
  ⟨proof⟩

setup-lifting type-definition-uint8

declare uint8.of-word-of [code abstype]

declare Quotient-uint8 [transfer-rule]

instantiation uint8 :: ⟨{comm-ring-1, semiring-modulo, equal, linorder}⟩
begin

lift-definition zero-uint8 :: uint8 is 0 ⟨proof⟩
lift-definition one-uint8 :: uint8 is 1 ⟨proof⟩
lift-definition plus-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨(+)⟩ ⟨proof⟩
lift-definition uminus-uint8 :: ⟨uint8 ⇒ uint8⟩ is uminus ⟨proof⟩
lift-definition minus-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨(-)⟩ ⟨proof⟩
lift-definition times-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨(*)⟩ ⟨proof⟩
```

```

lift-definition divide-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨(div)⟩ ⟨proof⟩
lift-definition modulo-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨(mod)⟩ ⟨proof⟩
lift-definition equal-uint8 :: ⟨uint8 ⇒ uint8 ⇒ bool⟩ is ⟨HOL.equal⟩ ⟨proof⟩
lift-definition less-eq-uint8 :: ⟨uint8 ⇒ uint8 ⇒ bool⟩ is ⟨(≤)⟩ ⟨proof⟩
lift-definition less-uint8 :: ⟨uint8 ⇒ uint8 ⇒ bool⟩ is ⟨(<)⟩ ⟨proof⟩

global-interpretation uint8: word-type-copy-ring Abs-uint8 Rep-uint8
  ⟨proof⟩

instance ⟨proof⟩

end

instantiation uint8 :: ring-bit-operations
begin

lift-definition bit-uint8 :: ⟨uint8 ⇒ nat ⇒ bool⟩ is bit ⟨proof⟩
lift-definition not-uint8 :: ⟨uint8 ⇒ uint8⟩ is ⟨Bit-Operations.not⟩ ⟨proof⟩
lift-definition and-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨Bit-Operations.and⟩
  ⟨proof⟩
lift-definition or-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨Bit-Operations.or⟩ ⟨proof⟩
lift-definition xor-uint8 :: ⟨uint8 ⇒ uint8 ⇒ uint8⟩ is ⟨Bit-Operations.xor⟩ ⟨proof⟩
lift-definition mask-uint8 :: ⟨nat ⇒ uint8⟩ is mask ⟨proof⟩
lift-definition push-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is push-bit ⟨proof⟩
lift-definition drop-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is drop-bit ⟨proof⟩
lift-definition signed-drop-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is signed-drop-bit
  ⟨proof⟩
lift-definition take-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is take-bit ⟨proof⟩
lift-definition set-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is Bit-Operations.set-bit
  ⟨proof⟩
lift-definition unset-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is unset-bit ⟨proof⟩
lift-definition flip-bit-uint8 :: ⟨nat ⇒ uint8 ⇒ uint8⟩ is flip-bit ⟨proof⟩

global-interpretation uint8: word-type-copy-bits Abs-uint8 Rep-uint8 signed-drop-bit-uint8
  ⟨proof⟩

instance
  ⟨proof⟩

end

lift-definition uint8-of-nat :: ⟨nat ⇒ uint8⟩
  is word-of-nat ⟨proof⟩

lift-definition nat-of-uint8 :: ⟨uint8 ⇒ nat⟩
  is unat ⟨proof⟩

lift-definition uint8-of-int :: ⟨int ⇒ uint8⟩
  is word-of-int ⟨proof⟩

```

```

lift-definition int-of-uint8 :: ⟨uint8 ⇒ int⟩
  is uint ⟨proof⟩

context
  includes integer.lifting
begin

lift-definition Uint8 :: ⟨integer ⇒ uint8⟩
  is word-of-int ⟨proof⟩

lift-definition integer-of-uint8 :: ⟨uint8 ⇒ integer⟩
  is uint ⟨proof⟩

end

global-interpretation uint8: word-type-copy-more Abs-uint8 Rep-uint8 signed-drop-bit-uint8
  uint8-of-nat nat-of-uint8 uint8-of-int int-of-uint8 Uint8 integer-of-uint8
  ⟨proof⟩

instantiation uint8 :: {size, msb, bit-comprehension}
begin

lift-definition size-uint8 :: ⟨uint8 ⇒ nat⟩ is size ⟨proof⟩

lift-definition msb-uint8 :: ⟨uint8 ⇒ bool⟩ is msb ⟨proof⟩

lift-definition set-bits-uint8 :: ⟨(nat ⇒ bool) ⇒ uint8⟩ is set-bits ⟨proof⟩
lift-definition set-bits-aux-uint8 :: ⟨(nat ⇒ bool) ⇒ nat ⇒ uint8 ⇒ uint8⟩ is
set-bits-aux ⟨proof⟩

global-interpretation uint8: word-type-copy-misc Abs-uint8 Rep-uint8 signed-drop-bit-uint8
  uint8-of-nat nat-of-uint8 uint8-of-int int-of-uint8 Uint8 integer-of-uint8 8 set-bits-aux-uint8
  ⟨proof⟩

instance ⟨proof⟩

end

```

7.2 Code setup

```

code-printing code-module Uint8 ↪ (SML)
  ⟨(* Test that words can handle numbers between 0 and 3 *)
  val - = if 3 <= Word.wordSize then () else raise (Fail (wordSize less than 3));

  structure Uint8 : sig
    val shiffl : Word8.word -> IntInf.int -> Word8.word
    val shiftr : Word8.word -> IntInf.int -> Word8.word
    val shiftr-signed : Word8.word -> IntInf.int -> Word8.word

```

```

    val test-bit : Word8.word -> IntInf.int -> bool
end = struct

fun shiftl x n =
    Word8.<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
    Word8.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
    Word8.~>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
    Word8.andb (x, Word8.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))) <>
    Word8.fromInt 0

end; (* struct Uint8 *)
code-reserved (SML) Uint8

code-printing code-module Uint8 -> (Haskell)
<module Uint8(Int8, Word8) where

    import Data.Int(Int8)
    import Data.Word(Word8)>
code-reserved (Haskell) Uint8

```

Scala provides only signed 8bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

```

code-printing code-module Uint8 -> (Scala)
<object Uint8 {

def less(x: Byte, y: Byte) : Boolean =
    x < 0 match {
        case true => y < 0 && x < y
        case false => y < 0 || x < y
    }

def less-eq(x: Byte, y: Byte) : Boolean =
    x < 0 match {
        case true => y < 0 && x <= y
        case false => y < 0 || x <= y
    }

def shiftl(x: Byte, n: BigInt) : Byte = (x << n.intValue).toByte

def shiftr(x: Byte, n: BigInt) : Byte = ((x & 255) >>> n.intValue).toByte

def shiftr-signed(x: Byte, n: BigInt) : Byte = (x >> n.intValue).toByte

```

```
def test-bit(x: Byte, n: BigInt) : Boolean =
  (x & (1 << n.intValue)) != 0
```

```
} /* object Uint8 */
code-reserved (Scala) Uint8
```

Avoid *Abs-uint8* in generated code, use *Rep-uint8'* instead. The symbolic implementations for `code_simp` use *Rep-uint8*.

The new destructor *Rep-uint8'* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint8* directly, because that makes `code_simp` loop.

If code generation raises `Match`, some equation probably contains *Rep-uint8* ([code abstract] equations for *uint8* may use *Rep-uint8* because these instances will be folded away.)

To convert *s* word values into *uint8*, use *Abs-uint8'*.

definition *Rep-uint8'* **where** [simp]: *Rep-uint8'* = *Rep-uint8*

lemma *Rep-uint8'-transfer* [transfer-rule]:
rel-fun cr-uint8 (=) ($\lambda x. x$) *Rep-uint8'*
 ⟨proof⟩

lemma *Rep-uint8'-code* [code]: *Rep-uint8'* *x* = (*BITS n. bit x n*)
 ⟨proof⟩

lift-definition *Abs-uint8'* :: *s* word \Rightarrow *uint8* **is** $\lambda x :: s$ word. *x* ⟨proof⟩

lemma *Abs-uint8'-code* [code]: *Abs-uint8'* *x* = *Uint8* (*integer-of-int* (*uint x*))
including *integer.lifting* ⟨proof⟩

declare [[code drop: *term-of-class.term-of* :: *uint8* \Rightarrow -]]

lemma *term-of-uint8-code* [code]:
defines *TR* \equiv *typerep.TypeRep* **and** *bit0* \equiv *STR "Numeral-Type.bit0"* **shows**
term-of-class.term-of *x* =
Code-Evaluation.App (*Code-Evaluation.Const* (*STR "Uint8.uint8.Abs-uint8"*)
 (*TR* (*STR "fun"*) [*TR* (*STR "Word.word"*) [*TR bit0* [*TR bit0* [*TR bit0* [*TR* (*STR*
"Numeral-Type.num1") []]]], *TR* (*STR "Uint8.uint8'"*) []]))
 (*term-of-class.term-of* (*Rep-uint8'* *x*))
 ⟨proof⟩

lemma *Uin8-code* [code]: *Rep-uint8* (*Uint8 i*) = *word-of-int* (*int-of-integer-symbolic*
i)
 ⟨proof⟩

code-printing type-constructor *uint8* \rightarrow
 (*SML*) *Word8.word* **and**
 (*Haskell*) *Uint8.Word8* **and**
 (*Scala*) *Byte*

```

| constant Uint8  $\rightarrow$ 
  (SML) Word8.fromLargeInt (IntInf.toLarge -) and
  (Haskell) (Prelude.fromInteger - :: Uint8.Word8) and
  (Haskell-Quickcheck) (Prelude.fromInteger (Prelude.toInteger -) :: Uint8.Word8)
and
  (Scala) -.byteValue
| constant 0 :: uint8  $\rightarrow$ 
  (SML) (Word8.fromInt 0) and
  (Haskell) (0 :: Uint8.Word8) and
  (Scala) 0.toByte
| constant 1 :: uint8  $\rightarrow$ 
  (SML) (Word8.fromInt 1) and
  (Haskell) (1 :: Uint8.Word8) and
  (Scala) 1.toByte
| constant plus :: uint8  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Word8.+ ((-, (-))) and
  (Haskell) infixl 6 + and
  (Scala) (- +/ -).toByte
| constant uminus :: uint8  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Word8.~ and
  (Haskell) negate and
  (Scala) (- -).toByte
| constant minus :: uint8  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Word8.- ((-, (-))) and
  (Haskell) infixl 6 - and
  (Scala) (- -/ -).toByte
| constant times :: uint8  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Word8.* ((-, (-))) and
  (Haskell) infixl 7 * and
  (Scala) (- */ -).toByte
| constant HOL.equal :: uint8  $\Rightarrow$  -  $\Rightarrow$  bool  $\rightarrow$ 
  (SML) !((- : Word8.word) = -) and
  (Haskell) infix 4 == and
  (Scala) infixl 5 ==
| class-instance uint8 :: equal  $\rightarrow$  (Haskell) -
| constant less-eq :: uint8  $\Rightarrow$  -  $\Rightarrow$  bool  $\rightarrow$ 
  (SML) Word8.<= ((-, (-))) and
  (Haskell) infix 4 <= and
  (Scala) Uint8.less'-eq
| constant less :: uint8  $\Rightarrow$  -  $\Rightarrow$  bool  $\rightarrow$ 
  (SML) Word8.< ((-, (-))) and
  (Haskell) infix 4 < and
  (Scala) Uint8.less
| constant Bit-Operations.not :: uint8  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Word8.notb and
  (Haskell) Data'-Bits.complement and
  (Scala) -.unary'~.toByte
| constant Bit-Operations.and :: uint8  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Word8.andb ((-, / (-))) and

```

```

(Haskell) infixl 7 Data-Bits..&. and
(Scala) (- & -).toByte
| constant Bit-Operations.or :: uint8 ⇒ - →
(SML) Word8.orb ((-),/ (-)) and
(Haskell) infixl 5 Data-Bits..|. and
(Scala) (- | -).toByte
| constant Bit-Operations.xor :: uint8 ⇒ - →
(SML) Word8.xorb ((-),/ (-)) and
(Haskell) Data'-Bits.xor and
(Scala) (- ^ -).toByte

```

definition *uint8-divmod* :: *uint8* ⇒ *uint8* ⇒ *uint8* × *uint8* **where**
uint8-divmod *x y* =
 (if *y* = 0 then (undefined ((*div*) :: *uint8* ⇒ -) *x* (0 :: *uint8*), undefined ((*mod*) ::
uint8 ⇒ -) *x* (0 :: *uint8*))
 else (*x div y*, *x mod y*)

definition *uint8-div* :: *uint8* ⇒ *uint8* ⇒ *uint8*
where *uint8-div* *x y* = *fst* (*uint8-divmod* *x y*)

definition *uint8-mod* :: *uint8* ⇒ *uint8* ⇒ *uint8*
where *uint8-mod* *x y* = *snd* (*uint8-divmod* *x y*)

lemma *div-uint8-code* [*code*]: *x div y* = (if *y* = 0 then 0 else *uint8-div* *x y*)
including *undefined-transfer* ⟨*proof*⟩

lemma *mod-uint8-code* [*code*]: *x mod y* = (if *y* = 0 then *x* else *uint8-mod* *x y*)
including *undefined-transfer* ⟨*proof*⟩

definition *uint8-sdiv* :: *uint8* ⇒ *uint8* ⇒ *uint8*
where
uint8-sdiv *x y* =
 (if *y* = 0 then undefined ((*div*) :: *uint8* ⇒ -) *x* (0 :: *uint8*)
 else *Abs-uint8* (*Rep-uint8* *x sdiv Rep-uint8* *y*)

definition *div0-uint8* :: *uint8* ⇒ *uint8*
where [*code del*]: *div0-uint8* *x* = undefined ((*div*) :: *uint8* ⇒ -) *x* (0 :: *uint8*)
declare [[*code abort: div0-uint8*]]

definition *mod0-uint8* :: *uint8* ⇒ *uint8*
where [*code del*]: *mod0-uint8* *x* = undefined ((*mod*) :: *uint8* ⇒ -) *x* (0 :: *uint8*)
declare [[*code abort: mod0-uint8*]]

lemma *uint8-divmod-code* [*code*]:
uint8-divmod *x y* =
 (if $0x80 \leq y$ then if $x < y$ then (0, *x*) else (1, *x* - *y*)
 else if *y* = 0 then (*div0-uint8* *x*, *mod0-uint8* *x*)
 else let *q* = *push-bit 1* (*uint8-sdiv* (*drop-bit 1* *x*) *y*);
 r = *x* - *q* * *y*

in if $r \geq y$ then $(q + 1, r - y)$ else (q, r)
including *undefined-transfer* *<proof>*

lemma *uint8-sdiv-code* [code]:

Rep-uint8 (uint8-sdiv x y) =
(if y = 0 then Rep-uint8 (undefined ((div) :: uint8 => -) x (0 :: uint8))
else Rep-uint8 x sdiv Rep-uint8 y)
<proof>

Note that we only need a translation for signed division, but not for the remainder because *uint8-divmod ?x ?y = (if 128 ≤ ?y then if ?x < ?y then (0, ?x) else (1, ?x - ?y) else if ?y = 0 then (div0-uint8 ?x, mod0-uint8 ?x) else let q = push-bit 1 (uint8-sdiv (drop-bit 1 ?x) ?y); r = ?x - q * ?y in if ?y ≤ r then (q + 1, r - ?y) else (q, r))* computes both with division only.

code-printing

constant *uint8-div* \rightarrow
(SML) Word8.div ((-), (-)) and
(Haskell) Prelude.div
| constant *uint8-mod* \rightarrow
(SML) Word8.mod ((-), (-)) and
(Haskell) Prelude.mod
| constant *uint8-divmod* \rightarrow
(Haskell) divmod
| constant *uint8-sdiv* \rightarrow
(Scala) (- '/' -).toByte

global-interpretation *uint8: word-type-copy-target-language Abs-uint8 Rep-uint8 signed-drop-bit-uint8*

uint8-of-nat nat-of-uint8 uint8-of-int int-of-uint8 Uint8 integer-of-uint8 8 set-bits-aux-uint8 8 7

defines *uint8-test-bit = uint8.test-bit*
and *uint8-shiffl = uint8.shiffl*
and *uint8-shiftr = uint8.shiftr*
and *uint8-sshiftr = uint8.sshiftr*
<proof>

code-printing constant *uint8-test-bit* \rightarrow

(SML) Uint8.test'-bit and
(Haskell) Data'-Bits.testBitBounded and
(Scala) Uint8.test'-bit and
(Eval) (fn w => fn i => if i < 0 orelse i >= 8 then raise (Fail argument to uint8'-test'-bit out of bounds) else Uint8.test'-bit w i)

code-printing constant *uint8-shiffl* \rightarrow

(SML) Uint8.shiffl and
(Haskell) Data'-Bits.shifflBounded and
(Scala) Uint8.shiffl and
(Eval) (fn w => fn i => if i < 0 orelse i >= 8 then raise (Fail argument to

uint8'-shiftr out of bounds) else *Uin8.shiftr w i*)

code-printing constant *uint8-shiftr* \rightarrow

(SML) *Uin8.shiftr* **and**

(Haskell) *Data'-Bits.shiftrBounded* **and**

(Scala) *Uin8.shiftr* **and**

(Eval) (fn w => fn i => if i < 0 orelse i >= 8 then raise (Fail argument to *uint8'-shiftr out of bounds*) else *Uin8.shiftr w i*)

code-printing constant *uint8-sshiftr* \rightarrow

(SML) *Uin8.shiftr'-signed* **and**

(Haskell)

(Prelude.fromInteger (Prelude.toInteger (Data'-Bits.shiftrBounded (Prelude.fromInteger (Prelude.toInteger -) :: Uin8.Int8) -)) :: Uin8.Word8) **and**

(Scala) *Uin8.shiftr'-signed* **and**

(Eval) (fn w => fn i => if i < 0 orelse i >= 8 then raise (Fail argument to *uint8'-sshiftr out of bounds*) else *Uin8.shiftr'-signed w i*)

context

includes *bit-operations-syntax*

begin

lemma *uint8-msb-test-bit*: *msb x* \longleftrightarrow *bit (x :: uint8) 7*

<proof>

lemma *msb-uint16-code* [code]: *msb x* \longleftrightarrow *uint8-test-bit x 7*

<proof>

lemma *uint8-of-int-code* [code]:

uint8-of-int i = Uin8 (integer-of-int i)

including *integer.lifting* *<proof>*

lemma *int-of-uint8-code* [code]:

int-of-uint8 x = int-of-integer (integer-of-uint8 x)

<proof>

lemma *uint8-of-nat-code* [code]:

uint8-of-nat = uint8-of-int o int

<proof>

lemma *nat-of-uint8-code* [code]:

nat-of-uint8 x = nat-of-integer (integer-of-uint8 x)

<proof> **including** *integer.lifting* *<proof>*

definition *integer-of-uint8-signed* :: *uint8* \Rightarrow *integer*

where

integer-of-uint8-signed n = (if bit n 7 then undefined integer-of-uint8 n else integer-of-uint8 n)

```

lemma integer-of-wint8-signed-code [code]:
  integer-of-wint8-signed n =
    (if bit n 7 then undefined integer-of-wint8 n else integer-of-int (uint (Rep-wint8'
n)))
  <proof>

```

```

lemma integer-of-wint8-code [code]:
  integer-of-wint8 n =
    (if bit n 7 then integer-of-wint8-signed (n AND 0x7F) OR 0x80 else integer-of-wint8-signed
n)
  <proof>
  including integer.lifting <proof>

```

end

code-printing

```

constant integer-of-wint8 →
  (SML) IntInf.fromLarge (Word8.toLargeInt -) and
  (Haskell) Prelude.toInteger
| constant integer-of-wint8-signed →
  (Scala) BigInt

```

7.3 Quickcheck setup

```

definition wint8-of-natural :: natural ⇒ wint8
where wint8-of-natural x ≡ Uint8 (integer-of-natural x)

```

```

instantiation wint8 :: {random, exhaustive, full-exhaustive} begin

```

```

definition random-wint8 ≡ qc-random-cnv wint8-of-natural

```

```

definition exhaustive-wint8 ≡ qc-exhaustive-cnv wint8-of-natural

```

```

definition full-exhaustive-wint8 ≡ qc-full-exhaustive-cnv wint8-of-natural

```

```

instance <proof>

```

end

```

instantiation wint8 :: narrowing begin

```

```

interpretation quickcheck-narrowing-samples

```

```

  λi. let x = Uint8 i in (x, 0xFF - x) 0

```

```

  Typerep.Typerep (STR "Uint8.wint8") [] <proof>

```

```

definition narrowing-wint8 d = qc-narrowing-drawn-from (narrowing-samples d)
d

```

```

declare [[code drop: partial-term-of :: wint8 itself ⇒ -]]

```

```

lemmas partial-term-of-wint8 [code] = partial-term-of-code

```

```

instance <proof>

```

end

end

Chapter 8

Unsigned words of default size

```
theory Uint
  imports
    HOL-Library.Code-Target-Bit-Shifts
    Uint-Common
    Code-Target-Word
    Code-Int-Integer-Conversion
begin
```

This theory provides access to words in the target languages of the code generator whose bit width is the default of the target language. To that end, the type *uint* models words of width *dflt-size*, but *dflt-size* is known only to be positive.

Usage restrictions: Default-size words (type *uint*) cannot be used for evaluation, because the results depend on the particular choice of word size in the target language and implementation. Symbolic evaluation has not yet been set up for *uint*.

The default size type

```
typedecl dflt-size
```

```
instantiation dflt-size :: typerep begin
```

```
definition typerep-class.typerep  $\equiv$   $\lambda-$  :: dflt-size itself. Typerep.TypeRep (STR  
"Uint.dflt-size") []
```

```
instance  $\langle$ proof $\rangle$ 
```

```
end
```

```
consts dflt-size-aux :: nat
```

```
specification (dflt-size-aux) dflt-size-aux-g0: dflt-size-aux > 0
```

```
 $\langle$ proof $\rangle$ 
```

```
hide-fact dflt-size-aux-def
```

```

instantiation dflt-size :: len begin
definition len-of-dflt-size (- :: dflt-size itself) ≡ dflt-size-aux
instance ⟨proof⟩
end

abbreviation dflt-size ≡ len-of (TYPE (dflt-size))

context includes integer.lifting begin
lift-definition dflt-size-integer :: integer is int dflt-size ⟨proof⟩
declare dflt-size-integer-def[code del]
  — The code generator will substitute a machine-dependent value for this constant

lemma dflt-size-by-int[code]: dflt-size = nat-of-integer dflt-size-integer
  ⟨proof⟩

lemma dflt-size[simp]:
  dflt-size > 0
  dflt-size ≥ Suc 0
  ¬ dflt-size < Suc 0
  ⟨proof⟩
end

```

8.1 Type definition and primitive operations

```

typedef uint = ⟨UNIV :: dflt-size word set⟩ ⟨proof⟩

global-interpretation uint: word-type-copy Abs-uint Rep-uint
  ⟨proof⟩

setup-lifting type-definition-uint

declare uint.of-word-of [code abstype]

declare Quotient-uint [transfer-rule]

instantiation uint :: ⟨{comm-ring-1, semiring-modulo, equal, linorder}⟩
begin

lift-definition zero-uint :: uint is 0 ⟨proof⟩
lift-definition one-uint :: uint is 1 ⟨proof⟩
lift-definition plus-uint :: ⟨uint ⇒ uint ⇒ uint⟩ is ⟨(+)⟩ ⟨proof⟩
lift-definition uminus-uint :: ⟨uint ⇒ uint⟩ is uminus ⟨proof⟩
lift-definition minus-uint :: ⟨uint ⇒ uint ⇒ uint⟩ is ⟨(-)⟩ ⟨proof⟩
lift-definition times-uint :: ⟨uint ⇒ uint ⇒ uint⟩ is ⟨(*)⟩ ⟨proof⟩
lift-definition divide-uint :: ⟨uint ⇒ uint ⇒ uint⟩ is ⟨(div)⟩ ⟨proof⟩
lift-definition modulo-uint :: ⟨uint ⇒ uint ⇒ uint⟩ is ⟨(mod)⟩ ⟨proof⟩
lift-definition equal-uint :: ⟨uint ⇒ uint ⇒ bool⟩ is ⟨HOL.equal⟩ ⟨proof⟩
lift-definition less-eq-uint :: ⟨uint ⇒ uint ⇒ bool⟩ is ⟨(≤)⟩ ⟨proof⟩
lift-definition less-uint :: ⟨uint ⇒ uint ⇒ bool⟩ is ⟨(<)⟩ ⟨proof⟩

```

global-interpretation *uint*: *word-type-copy-ring Abs-uint Rep-uint*
 ⟨*proof*⟩

instance ⟨*proof*⟩

end

instantiation *uint* :: *ring-bit-operations*
begin

lift-definition *bit-uint* :: ⟨*uint* ⇒ *nat* ⇒ *bool*⟩ **is** *bit* ⟨*proof*⟩

lift-definition *not-uint* :: ⟨*uint* ⇒ *uint*⟩ **is** ⟨*Bit-Operations.not*⟩ ⟨*proof*⟩

lift-definition *and-uint* :: ⟨*uint* ⇒ *uint* ⇒ *uint*⟩ **is** ⟨*Bit-Operations.and*⟩ ⟨*proof*⟩

lift-definition *or-uint* :: ⟨*uint* ⇒ *uint* ⇒ *uint*⟩ **is** ⟨*Bit-Operations.or*⟩ ⟨*proof*⟩

lift-definition *xor-uint* :: ⟨*uint* ⇒ *uint* ⇒ *uint*⟩ **is** ⟨*Bit-Operations.xor*⟩ ⟨*proof*⟩

lift-definition *mask-uint* :: ⟨*nat* ⇒ *uint*⟩ **is** *mask* ⟨*proof*⟩

lift-definition *push-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *push-bit* ⟨*proof*⟩

lift-definition *drop-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *drop-bit* ⟨*proof*⟩

lift-definition *signed-drop-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *signed-drop-bit* ⟨*proof*⟩

lift-definition *take-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *take-bit* ⟨*proof*⟩

lift-definition *set-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *Bit-Operations.set-bit* ⟨*proof*⟩

lift-definition *unset-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *unset-bit* ⟨*proof*⟩

lift-definition *flip-bit-uint* :: ⟨*nat* ⇒ *uint* ⇒ *uint*⟩ **is** *flip-bit* ⟨*proof*⟩

global-interpretation *uint*: *word-type-copy-bits Abs-uint Rep-uint signed-drop-bit-uint*
 ⟨*proof*⟩

instance
 ⟨*proof*⟩

end

lift-definition *uint-of-nat* :: ⟨*nat* ⇒ *uint*⟩
is *word-of-nat* ⟨*proof*⟩

lift-definition *nat-of-uint* :: ⟨*uint* ⇒ *nat*⟩
is *unat* ⟨*proof*⟩

lift-definition *uint-of-int* :: ⟨*int* ⇒ *uint*⟩
is *word-of-int* ⟨*proof*⟩

lift-definition *int-of-uint* :: ⟨*uint* ⇒ *int*⟩
is *uint* ⟨*proof*⟩

context
includes *integer.lifting*
begin

lift-definition *Uint* :: $\langle \text{integer} \Rightarrow \text{uint} \rangle$
is *word-of-int* $\langle \text{proof} \rangle$

lift-definition *integer-of-uint* :: $\langle \text{uint} \Rightarrow \text{integer} \rangle$
is *uint* $\langle \text{proof} \rangle$

end

global-interpretation *uint*: *word-type-copy-more Abs-uint Rep-uint signed-drop-bit-uint*
uint-of-nat nat-of-uint uint-of-int int-of-uint Uint integer-of-uint
 $\langle \text{proof} \rangle$

instantiation *uint* :: $\{ \text{size}, \text{msb}, \text{bit-comprehension} \}$
begin

lift-definition *size-uint* :: $\langle \text{uint} \Rightarrow \text{nat} \rangle$ **is** *size* $\langle \text{proof} \rangle$

lift-definition *msb-uint* :: $\langle \text{uint} \Rightarrow \text{bool} \rangle$ **is** *msb* $\langle \text{proof} \rangle$

lift-definition *set-bits-uint* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{uint} \rangle$ **is** *set-bits* $\langle \text{proof} \rangle$

lift-definition *set-bits-aux-uint* :: $\langle (\text{nat} \Rightarrow \text{bool}) \Rightarrow \text{nat} \Rightarrow \text{uint} \Rightarrow \text{uint} \rangle$ **is** *set-bits-aux*
 $\langle \text{proof} \rangle$

global-interpretation *uint*: *word-type-copy-misc Abs-uint Rep-uint signed-drop-bit-uint*
uint-of-nat nat-of-uint uint-of-int int-of-uint Uint integer-of-uint dflt-size set-bits-aux-uint
 $\langle \text{proof} \rangle$

instance $\langle \text{proof} \rangle$

end

8.2 Code setup

code-printing code-module *Uint* \rightarrow (*SML*)

\langle

structure Uint : *sig*

val shiftl : *Word.word* \rightarrow *IntInf.int* \rightarrow *Word.word*

val shiftr : *Word.word* \rightarrow *IntInf.int* \rightarrow *Word.word*

val shiftr-signed : *Word.word* \rightarrow *IntInf.int* \rightarrow *Word.word*

val test-bit : *Word.word* \rightarrow *IntInf.int* \rightarrow *bool*

end = *struct*

fun shiftl *x n* =

Word.<< (*x*, *Word.fromLargeInt* (*IntInf.toLarge n*))

fun shiftr *x n* =

Word.>> (*x*, *Word.fromLargeInt* (*IntInf.toLarge n*))

fun shiftr-signed *x n* =

```

Word.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
  Word.andb (x, Word.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))) <>
  Word.fromInt 0

end; (* struct Uint *)
code-reserved (SML) Uint

code-printing code-module Uint  $\rightarrow$  (Haskell)
<module Uint(Int, Word, dflt-size) where

  import qualified Prelude
  import Data.Int(Int)
  import Data.Word(Word)
  import qualified Data.Bits

  dflt-size :: Prelude.Integer
  dflt-size = Prelude.toInteger (bitSize-aux (0::Word)) where
    bitSize-aux :: (Data.Bits.Bits a, Prelude.Bounded a) => a -> Int
    bitSize-aux = Data.Bits.bitSize
  and (Haskell-Quickcheck)
  <module Uint(Int, Word, dflt-size) where

    import qualified Prelude
    import Data.Int(Int)
    import Data.Word(Word)
    import qualified Data.Bits

    dflt-size :: Prelude.Int
    dflt-size = bitSize-aux (0::Word) where
      bitSize-aux :: (Data.Bits.Bits a, Prelude.Bounded a) => a -> Int
      bitSize-aux = Data.Bits.bitSize
  >
code-reserved (Haskell) Uint dflt-size

```

OCaml and Scala provide only signed bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

```

code-printing code-module Uint  $\rightarrow$  (OCaml)
<module Uint : sig
  type t = int
  val dflt-size : Z.t
  val less : t -> t -> bool
  val less-eq : t -> t -> bool
  val shiftl : t -> Z.t -> t
  val shiftr : t -> Z.t -> t
  val shiftr-signed : t -> Z.t -> t
  val test-bit : t -> Z.t -> bool
  val int-mask : int

```

```

    val int32-mask : int32
    val int64-mask : int64
end = struct

type t = int

let dflt-size = Z.of-int Sys.int-size;;

(* negative numbers have their highest bit set,
   so they are greater than positive ones *)
let less x y =
  if x < 0 then
    y < 0 && x < y
  else y < 0 || x < y;;

let less-eq x y =
  if x < 0 then
    y < 0 && x <= y
  else y < 0 || x <= y;;

let shiftl x n = x lsl (Z.to-int n);;

let shiftr x n = x lsr (Z.to-int n);;

let shiftr-signed x n = x asr (Z.to-int n);;

let test-bit x n = x land (1 lsl (Z.to-int n)) <> 0;;

let int-mask =
  if Sys.int-size < 32 then lnot 0 else 0xFFFFFFFF;;

let int32-mask =
  if Sys.int-size < 32 then Int32.pred (Int32.shift-left Int32.one Sys.int-size)
  else Int32.of-string 0xFFFFFFFF;;

let int64-mask =
  if Sys.int-size < 64 then Int64.pred (Int64.shift-left Int64.one Sys.int-size)
  else Int64.of-string 0xFFFFFFFFFFFFFFFF;;

end;; (*struct Uint*)
code-reserved (OCaml) Uint

code-printing code-module Uint → (Scala)
⟨object Uint {
def dflt-size : BigInt = BigInt(32)

def less(x: Int, y: Int) : Boolean =
  x < 0 match {
    case true => y < 0 && x < y

```

```

    case false => y < 0 || x < y
  }

def less-eq(x: Int, y: Int) : Boolean =
  x < 0 match {
    case true => y < 0 && x <= y
    case false => y < 0 || x <= y
  }

def shiftl(x: Int, n: BigInt) : Int = x << n.intValue

def shiftr(x: Int, n: BigInt) : Int = x >>> n.intValue

def shiftr-signed(x: Int, n: BigInt) : Int = x >> n.intValue

def test-bit(x: Int, n: BigInt) : Boolean =
  (x & (1 << n.intValue)) != 0

} /* object Uint */
code-reserved (Scala) Uint

```

OCaml's conversion from `Big_int` to `int` demands that the value fits into a signed integer. The following justifies the implementation.

context

includes *integer.lifting* and *bit-operations-syntax*
begin

definition *wivs-mask* :: *int* **where** *wivs-mask* = $2^{\text{dft-size} - 1}$

lift-definition *wivs-mask-integer* :: *integer* **is** *wivs-mask* *<proof>*

lemma [*code*]: *wivs-mask-integer* = $2^{\text{dft-size} - 1}$
<proof>

definition *wivs-shift* :: *int* **where** *wivs-shift* = $2^{\text{dft-size}}$

lift-definition *wivs-shift-integer* :: *integer* **is** *wivs-shift* *<proof>*

lemma [*code*]: *wivs-shift-integer* = $2^{\text{dft-size}}$
<proof>

definition *wivs-index* :: *nat* **where** *wivs-index* == $\text{dft-size} - 1$

lift-definition *wivs-index-integer* :: *integer* **is** *int wivs-index* *<proof>*

lemma *wivs-index-integer-code*[*code*]: *wivs-index-integer* = $\text{dft-size-integer} - 1$
<proof>

definition *wivs-overflow* :: *int* **where** *wivs-overflow* == $2^{(\text{dft-size} - 1)}$

lift-definition *wivs-overflow-integer* :: *integer* **is** *wivs-overflow* *<proof>*

lemma [*code*]: *wivs-overflow-integer* = $2^{(\text{dft-size} - 1)}$
<proof>

definition *wivs-least* :: *int* **where** *wivs-least* == $- \text{wivs-overflow}$

lift-definition *wivs-least-integer* :: *integer* **is** *wivs-least* *<proof>*

lemma [code]: $wivs\text{-}least\text{-}integer = -(2^{dflt\text{-}size - 1})$
 ⟨proof⟩

definition $Uint\text{-}signed :: integer \Rightarrow uint$ **where**

$Uint\text{-}signed\ i = (if\ i < wivs\text{-}least\text{-}integer \vee wivs\text{-}overflow\text{-}integer \leq i\ then\ unde\text{-}fined\ Uint\ i\ else\ Uint\ i)$

lemma $Uint\text{-}code$ [code]:

$Uint\ i =$
 (let $i' = i$ AND $wivs\text{-}mask\text{-}integer\ in$
 if $bit\ i'\ wivs\text{-}index$ then $Uint\text{-}signed\ (i' - wivs\text{-}shift\text{-}integer)$ else $Uint\text{-}signed\ i'$)
including $undefined\text{-}transfer$
 ⟨proof⟩

lemma $Uint\text{-}signed\text{-}code$ [code]:

$Rep\text{-}uint\ (Uint\text{-}signed\ i) =$
 (if $i < wivs\text{-}least\text{-}integer \vee i \geq wivs\text{-}overflow\text{-}integer$ then $Rep\text{-}uint\ (undefined\ Uint\ i)$ else $word\text{-}of\text{-}int\ (int\text{-}of\text{-}integer\text{-}symbolic\ i)$)
 ⟨proof⟩

end

Avoid $Abs\text{-}uint$ in generated code, use $Rep\text{-}uint'$ instead. The symbolic implementations for $code_simp$ use $Rep\text{-}uint$.

The new destructor $Rep\text{-}uint'$ is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement $Rep\text{-}uint$ directly, because that makes $code_simp$ loop.

If code generation raises Match, some equation probably contains $Rep\text{-}uint$ ([code abstract] equations for $uint$ may use $Rep\text{-}uint$ because these instances will be folded away.)

definition $Rep\text{-}uint'$ **where** [simp]: $Rep\text{-}uint' = Rep\text{-}uint$

lemma $Rep\text{-}uint'\text{-}code$ [code]: $Rep\text{-}uint'\ x = (BITS\ n.\ bit\ x\ n)$
 ⟨proof⟩

lift-definition $Abs\text{-}uint' :: dflt\text{-}size\ word \Rightarrow uint$ **is** $\lambda x :: dflt\text{-}size\ word.\ x$ ⟨proof⟩

lemma $Abs\text{-}uint'\text{-}code$ [code]:

$Abs\text{-}uint'\ x = Uint\ (integer\text{-}of\text{-}int\ (uint\ x))$

including $integer.\text{lifting}$ ⟨proof⟩

declare [[code drop: $term\text{-}of\text{-}class.\text{term}\text{-}of :: uint \Rightarrow -$]]

lemma $term\text{-}of\text{-}uint\text{-}code$ [code]:

defines $TR \equiv typerep.\text{Typerep}$ **and** $bit0 \equiv STR\ "Numeral\text{-}Type.\text{bit}0"$

shows

$term\text{-}of\text{-}class.\text{term}\text{-}of\ x =$

$Code\text{-}Evaluation.\text{App}\ (Code\text{-}Evaluation.\text{Const}\ (STR\ "Uint.\text{uint}.\text{Abs}\text{-}\text{uint}"))\ (TR\ (STR\ "fun"))\ [TR\ (STR\ "Word.\text{word}"))\ [TR\ (STR\ "Uint.\text{dflt}\text{-}\text{size}"))\ []],\ TR\ (STR$

```
"Uint.uint'" []])
  (term-of-class.term-of (Rep-uint' x))
  ⟨proof⟩
```

Important: We must prevent the reflection oracle (`eval-tac`) to use our machine-dependent type.

code-printing

```
type-constructor uint →
  (SML) Word.word and
  (Haskell) Uint.Word and
  (OCaml) Uint.t and
  (Scala) Int and
  (Eval) *** Error: Machine dependent type *** and
  (Quickcheck) Word.word
| constant dflt-size-integer →
  (SML) (IntInf.fromLarge (Int.toLarge Word.wordSize)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.wordSize and
  (Haskell) Uint.dflt'-size and
  (OCaml) Uint.dflt'-size and
  (Scala) Uint.dflt'-size
| constant Uint →
  (SML) Word.fromLargeInt (IntInf.toLarge -) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.fromInt and
  (Haskell) (Prelude.fromInteger - :: Uint.Word) and
  (Haskell-Quickcheck) (Prelude.fromInteger (Prelude.toInteger -) :: Uint.Word)
and
  (Scala) -.intValue
| constant Uint-signed →
  (OCaml) Z.to'-int
| constant 0 :: uint →
  (SML) (Word.fromInt 0) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) (Word.fromInt 0) and
  (Haskell) (0 :: Uint.Word) and
  (OCaml) 0 and
  (Scala) 0
| constant 1 :: uint →
  (SML) (Word.fromInt 1) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) (Word.fromInt 1) and
  (Haskell) (1 :: Uint.Word) and
  (OCaml) 1 and
  (Scala) 1
| constant plus :: uint ⇒ - →
  (SML) Word.+ ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.+ ((-), (-)) and
```

```

(Haskell) infixl 6 + and
(OCaml) Pervasives.(+) and
(Scala) infixl 7 +
| constant uminus :: uint ⇒ - →
(SML) Word.~ and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.~ and
(Haskell) negate and
(OCaml) Pervasives.(~ -) and
(Scala) !(- -)
| constant minus :: uint ⇒ - →
(SML) Word.- ((-), (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.- ((-), (-)) and
(Haskell) infixl 6 - and
(OCaml) Pervasives.(-) and
(Scala) infixl 7 -
| constant times :: uint ⇒ - ⇒ - →
(SML) Word.* ((-), (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.* ((-), (-)) and
(Haskell) infixl 7 * and
(OCaml) Pervasives.( * ) and
(Scala) infixl 8 *
| constant HOL.equal :: uint ⇒ - ⇒ bool →
(SML) !((- : Word.word) = -) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) !((- : Word.word) = -) and
(Haskell) infix 4 == and
(OCaml) (Pervasives.(=):Uint.t -> Uint.t -> bool) and
(Scala) infixl 5 ==
| class-instance uint :: equal →
(Haskell) -
| constant less-eq :: uint ⇒ - ⇒ bool →
(SML) Word.<= ((-), (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.<= ((-), (-)) and
(Haskell) infix 4 <= and
(OCaml) Uint.less'-eq and
(Scala) Uint.less'-eq
| constant less :: uint ⇒ - ⇒ bool →
(SML) Word.< ((-), (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.< ((-), (-)) and
(Haskell) infix 4 < and
(OCaml) Uint.less and
(Scala) Uint.less
| constant Bit-Operations.not :: uint ⇒ - →
(SML) Word.notb and

```

```

(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.notb and
(Haskell) Data'-Bits.complement and
(OCaml) Pervasives.lnot and
(Scala) -.unary'~
| constant Bit-Operations.and :: uint ⇒ - →
(SML) Word.andb ((-),/ (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.andb ((-),/ (-)) and
(Haskell) infixl 7 Data-Bits.&. and
(OCaml) Pervasives.(land) and
(Scala) infixl 3 &
| constant Bit-Operations.or :: uint ⇒ - →
(SML) Word.orb ((-),/ (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.orb ((-),/ (-)) and
(Haskell) infixl 5 Data-Bits.|. and
(OCaml) Pervasives.(lor) and
(Scala) infixl 1 |
| constant Bit-Operations.xor :: uint ⇒ - →
(SML) Word.xorb ((-),/ (-)) and
(Eval) (raise (Fail Machine dependent code)) and
(Quickcheck) Word.xorb ((-),/ (-)) and
(Haskell) Data'-Bits.xor and
(OCaml) Pervasives.(lxor) and
(Scala) infixl 2 ^

definition uint-divmod :: uint ⇒ uint ⇒ uint × uint where
  uint-divmod x y =
    (if y = 0 then (undefined ((div) :: uint ⇒ -) x (0 :: uint), undefined ((mod) ::
uint ⇒ -) x (0 :: uint))
    else (x div y, x mod y))

definition uint-div :: uint ⇒ uint ⇒ uint
where uint-div x y = fst (uint-divmod x y)

definition uint-mod :: uint ⇒ uint ⇒ uint
where uint-mod x y = snd (uint-divmod x y)

lemma div-uint-code [code]: x div y = (if y = 0 then 0 else uint-div x y)
including undefined-transfer ⟨proof⟩

lemma mod-uint-code [code]: x mod y = (if y = 0 then x else uint-mod x y)
including undefined-transfer ⟨proof⟩

definition uint-sdiv :: uint ⇒ uint ⇒ uint
where [code del]:
  uint-sdiv x y =
    (if y = 0 then undefined ((div) :: uint ⇒ -) x (0 :: uint)

```

else Abs-uint (Rep-uint x sdiv Rep-uint y)

definition *div0-uint* :: *uint* \Rightarrow *uint*

where [*code del*]: *div0-uint* *x* = *undefined* ((*div*) :: *uint* \Rightarrow -) *x* (*0* :: *uint*)

declare [[*code abort*: *div0-uint*]]

definition *mod0-uint* :: *uint* \Rightarrow *uint*

where [*code del*]: *mod0-uint* *x* = *undefined* ((*mod*) :: *uint* \Rightarrow -) *x* (*0* :: *uint*)

declare [[*code abort*: *mod0-uint*]]

definition *wivs-overflow-uint* :: *uint*

where *wivs-overflow-uint* \equiv *push-bit* (*dflt-size* - 1) 1

lemma *Rep-uint-wivs-overflow-uint-eq*:

\langle *Rep-uint wivs-overflow-uint* = 2^{\wedge} (*dflt-size* - *Suc* 0) \rangle

\langle *proof* \rangle

lemma *wivs-overflow-uint-greater-eq-0*:

\langle *wivs-overflow-uint* > 0 \rangle

\langle *proof* \rangle

lemma *uint-divmod-code* [*code*]:

uint-divmod *x* *y* =

(*if wivs-overflow-uint* \leq *y* then *if* *x* < *y* then (*0*, *x*) else (*1*, *x* - *y*)

else *if* *y* = 0 then (*div0-uint* *x*, *mod0-uint* *x*)

else let *q* = *push-bit* 1 (*uint-sdiv* (*drop-bit* 1 *x*) *y*);

r = *x* - *q* * *y*

in *if* *r* \geq *y* then (*q* + 1, *r* - *y*) else (*q*, *r*)

\langle *proof* \rangle

including *undefined-transfer*

\langle *proof* \rangle

lemma *uint-sdiv-code* [*code*]:

Rep-uint (*uint-sdiv* *x* *y*) =

(*if* *y* = 0 then *Rep-uint* (*undefined* ((*div*) :: *uint* \Rightarrow -) *x* (*0* :: *uint*))

else *Rep-uint* *x* *sdiv* *Rep-uint* *y*)

\langle *proof* \rangle

Note that we only need a translation for signed division, but not for the remainder because *uint-divmod* ?*x* ?*y* = (*if wivs-overflow-uint* \leq ?*y* then *if* ?*x* < ?*y* then (*0*, ?*x*) else (*1*, ?*x* - ?*y*) else *if* ?*y* = 0 then (*div0-uint* ?*x*, *mod0-uint* ?*x*) else let *q* = *push-bit* 1 (*uint-sdiv* (*drop-bit* 1 ?*x*) ?*y*); *r* = ?*x* - *q* * ?*y* in *if* ?*y* \leq *r* then (*q* + 1, *r* - ?*y*) else (*q*, *r*)) computes both with division only.

code-printing

constant *uint-div* \rightarrow

(*SML*) *Word.div* ((-), (-)) **and**

(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**

(*Quickcheck*) *Word.div* ((-), (-)) **and**

```

(Haskell) Prelude.div
| constant uint-mod  $\rightarrow$ 
  (SML) Word.mod ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.mod ((-), (-)) and
  (Haskell) Prelude.mod
| constant uint-divmod  $\rightarrow$ 
  (Haskell) divmod
| constant uint-sdiv  $\rightarrow$ 
  (OCaml) Pervasives.(/) and
  (Scala) - / -

```

global-interpretation *uint*: *word-type-copy-target-language Abs-uint Rep-uint signed-drop-bit-uint*
uint-of-nat nat-of-uint uint-of-int int-of-uint Uint integer-of-uint dflt-size set-bits-aux-uint

```

<of-nat dflt-size> wivs-index
defines uint-test-bit = uint.test-bit
  and uint-shiffl = uint.shiffl
  and uint-shiftr = uint.shiftr
  and uint-sshiftr = uint.sshiftr
<proof>

```

```

code-printing constant uint-test-bit  $\rightarrow$ 
  (SML) Uint.test'-bit and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Uint.test'-bit and
  (Haskell) Data'-Bits.testBitBounded and
  (OCaml) Uint.test'-bit and
  (Scala) Uint.test'-bit

```

```

code-printing constant uint-shiffl  $\rightarrow$ 
  (SML) Uint.shiffl and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Uint.shiffl and
  (Haskell) Data'-Bits.shifflBounded and
  (OCaml) Uint.shiffl and
  (Scala) Uint.shiffl

```

```

code-printing constant uint-shiftr  $\rightarrow$ 
  (SML) Uint.shiftr and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Uint.shiftr and
  (Haskell) Data'-Bits.shiftrBounded and
  (OCaml) Uint.shiftr and
  (Scala) Uint.shiftr

```

```

code-printing constant uint-sshiftr  $\rightarrow$ 
  (SML) Uint.shiftr'-signed and
  (Eval) (raise (Fail Machine dependent code)) and

```

```

(Quickcheck) Uint.shiftr'-signed and
(Haskell)
  (Prelude.fromInteger (Prelude.toInteger (Data'-Bits.shiftrBounded (Prelude.fromInteger
(Prelude.toInteger -) :: Uint.Int) -)) :: Uint.Word) and
(OCaml) Uint.shiftr'-signed and
(Scala) Uint.shiftr'-signed

```

```

lemma uint-msb-test-bit: msb  $x \longleftrightarrow$  bit ( $x ::$  uint) wivs-index
  <proof>

```

```

lemma msb-uint-code [code]: msb  $x \longleftrightarrow$  uint-test-bit  $x$  wivs-index-integer
  <proof>

```

```

lemma uint-of-int-code [code]: uint-of-int  $i =$  (BITS  $n$ . bit  $i$   $n$ )
  <proof>

```

8.3 Quickcheck setup

```

definition uint-of-natural :: natural  $\Rightarrow$  uint
where uint-of-natural  $x \equiv$  Uint (integer-of-natural  $x$ )

```

```

instantiation uint :: {random, exhaustive, full-exhaustive} begin
definition random-uint  $\equiv$  qc-random-cnv uint-of-natural
definition exhaustive-uint  $\equiv$  qc-exhaustive-cnv uint-of-natural
definition full-exhaustive-uint  $\equiv$  qc-full-exhaustive-cnv uint-of-natural
instance <proof>
end

```

```

instantiation uint :: narrowing begin

```

```

interpretation quickcheck-narrowing-samples
   $\lambda i. (Uint\ i, Uint\ (-\ i))\ 0$ 
  Typerep.Typerep (STR "Uint.uint") [] <proof>

```

```

definition narrowing-uint  $d =$  qc-narrowing-drawn-from (narrowing-samples  $d$ )  $d$ 
declare [[code drop: partial-term-of :: uint itself  $\Rightarrow$  -]]
lemmas partial-term-of-uint [code] = partial-term-of-code

```

```

instance <proof>
end

```

```

find-consts name: wivs

```

```

end

```

Chapter 9

Conversions between unsigned words and between char

theory *Native-Cast*

imports

Uint8

Uint16

Uint32

Uint64

begin

Auxiliary stuff

lemma *integer-of-char-char-of-integer* [*simp*]:
 $integer\text{-of-char } (char\text{-of-integer } x) = x \bmod 256$
<proof>

lemma *char-of-integer-integer-of-char* [*simp*]:
 $char\text{-of-integer } (integer\text{-of-char } x) = x$
<proof>

lemma *int-lt-numeral* [*simp*]: $int\ x < numeral\ n \iff x < numeral\ n$
<proof>

lemma *int-of-integer-ge-0*: $0 \leq int\text{-of-integer } x \iff 0 \leq x$
including *integer.lifting* *<proof>*

lemma *integer-of-char-ge-0* [*simp*]: $0 \leq integer\text{-of-char } x$
including *integer.lifting* *<proof>*

9.1 Conversion between native words

lift-definition *uint8-of-uint16* :: $uint16 \Rightarrow uint8$ **is** *ucast* *<proof>*

lift-definition *wint8-of-wint32* :: *wint32* ⇒ *wint8* **is** *ucast* ⟨*proof*⟩

lift-definition *wint8-of-wint64* :: *wint64* ⇒ *wint8* **is** *ucast* ⟨*proof*⟩

lift-definition *wint16-of-wint8* :: *wint8* ⇒ *wint16* **is** *ucast* ⟨*proof*⟩

lift-definition *wint16-of-wint32* :: *wint32* ⇒ *wint16* **is** *ucast* ⟨*proof*⟩

lift-definition *wint16-of-wint64* :: *wint64* ⇒ *wint16* **is** *ucast* ⟨*proof*⟩

lift-definition *wint32-of-wint8* :: *wint8* ⇒ *wint32* **is** *ucast* ⟨*proof*⟩

lift-definition *wint32-of-wint16* :: *wint16* ⇒ *wint32* **is** *ucast* ⟨*proof*⟩

lift-definition *wint32-of-wint64* :: *wint64* ⇒ *wint32* **is** *ucast* ⟨*proof*⟩

lift-definition *wint64-of-wint8* :: *wint8* ⇒ *wint64* **is** *ucast* ⟨*proof*⟩

lift-definition *wint64-of-wint16* :: *wint16* ⇒ *wint64* **is** *ucast* ⟨*proof*⟩

lift-definition *wint64-of-wint32* :: *wint32* ⇒ *wint64* **is** *ucast* ⟨*proof*⟩

context

begin

qualified definition *mask* :: *integer*

where ⟨*mask* = (*0xFFFFFFFF* :: *integer*)⟩

end

code-printing

constant *wint8-of-wint16* →

(*SML-word*) *Word8.fromLarge* (*Word16.toLarge* -) **and**

(*Haskell*) (*Prelude.fromIntegral* - :: *Uint8.Word8*) **and**

(*Scala*) -.*toByte*

| **constant** *wint8-of-wint32* →

(*SML*) *Word8.fromLarge* (*Word32.toLarge* -) **and**

(*Haskell*) (*Prelude.fromIntegral* - :: *Uint8.Word8*) **and**

(*Scala*) -.*toByte*

| **constant** *wint8-of-wint64* →

(*SML*) *Word8.fromLarge* (*Uint64.toLarge* -) **and**

(*Haskell*) (*Prelude.fromIntegral* - :: *Uint8.Word8*) **and**

(*Scala*) -.*toByte*

| **constant** *wint16-of-wint8* →

(*SML-word*) *Word16.fromLarge* (*Word8.toLarge* -) **and**

(*Haskell*) (*Prelude.fromIntegral* - :: *Uint16.Word16*) **and**

(*Scala*) ((-).*toInt* & *0xFF*).*toChar*

| **constant** *wint16-of-wint32* →

(*SML-word*) *Word16.fromLarge* (*Word32.toLarge* -) **and**

(*Haskell*) (*Prelude.fromIntegral* - :: *Uint16.Word16*) **and**

(*Scala*) -.*toChar*

| **constant** *wint16-of-wint64* →

(*SML-word*) *Word16.fromLarge* (*Uint64.toLarge* -) **and**

(*Haskell*) (*Prelude.fromIntegral* - :: *Uint16.Word16*) **and**

(*Scala*) -.*toChar*

| **constant** *wint32-of-wint8* →

```

(SML) Word32.fromLarge (Word8.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
(Scala) ((-).toInt & 0xFF)
| constant uint32-of-uint16  $\rightarrow$ 
(SML-word) Word32.fromLarge (Word16.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
(Scala) (-).toInt
| constant uint32-of-uint64  $\rightarrow$ 
(SML-word) Word32.fromLarge (Uint64.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
(Scala) (-).toInt and
(OCaml) Int64.to'-int32
| constant uint64-of-uint8  $\rightarrow$ 
(SML-word) Word64.fromLarge (Word8.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
(Scala) ((-).toLong & 0xFF)
| constant uint64-of-uint16  $\rightarrow$ 
(SML-word) Word64.fromLarge (Word16.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
(Scala) -.toLong
| constant uint64-of-uint32  $\rightarrow$ 
(SML-word) Word64.fromLarge (Word32.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
(Scala) ((-).toLong & 0xFFFFFFFFL) and
(OCaml) Int64.logand (Int64.of'-int32 -) (Int64.of'-string 4294967295)

```

Use *Abs-uint8'* etc. instead of *Rep-uint8* in code equations for conversion functions to avoid exceptions during code generation when the target language provides only some of the uint types.

```

lemma uint8-of-uint16-code [code]:
  uint8-of-uint16 x = Abs-uint8' (ucast (Rep-uint16' x))
<proof>

```

```

lemma uint8-of-uint32-code [code]:
  uint8-of-uint32 x = Abs-uint8' (ucast (Rep-uint32' x))
<proof>

```

```

lemma uint8-of-uint64-code [code]:
  uint8-of-uint64 x = Abs-uint8' (ucast (Rep-uint64' x))
<proof>

```

```

lemma uint16-of-uint8-code [code]:
  uint16-of-uint8 x = Abs-uint16' (ucast (Rep-uint8' x))
<proof>

```

```

lemma uint16-of-uint32-code [code]:
  uint16-of-uint32 x = Abs-uint16' (ucast (Rep-uint32' x))
<proof>

```

lemma *uint16-of-uint64-code* [code]:
 $\text{uint16-of-uint64 } x = \text{Abs-uint16}' (\text{ucast } (\text{Rep-uint64}' x))$
 ⟨proof⟩

lemma *uint32-of-uint8-code* [code]:
 $\text{uint32-of-uint8 } x = \text{Abs-uint32}' (\text{ucast } (\text{Rep-uint8}' x))$
 ⟨proof⟩

lemma *uint32-of-uint16-code* [code]:
 $\text{uint32-of-uint16 } x = \text{Abs-uint32}' (\text{ucast } (\text{Rep-uint16}' x))$
 ⟨proof⟩

lemma *uint32-of-uint64-code* [code]:
 $\text{uint32-of-uint64 } x = \text{Abs-uint32}' (\text{ucast } (\text{Rep-uint64}' x))$
 ⟨proof⟩

lemma *uint64-of-uint8-code* [code]:
 $\text{uint64-of-uint8 } x = \text{Abs-uint64}' (\text{ucast } (\text{Rep-uint8}' x))$
 ⟨proof⟩

lemma *uint64-of-uint16-code* [code]:
 $\text{uint64-of-uint16 } x = \text{Abs-uint64}' (\text{ucast } (\text{Rep-uint16}' x))$
 ⟨proof⟩

lemma *uint64-of-uint32-code* [code]:
 $\text{uint64-of-uint32 } x = \text{Abs-uint64}' (\text{ucast } (\text{Rep-uint32}' x))$
 ⟨proof⟩

end

theory *Native-Cast-Uint* **imports**

Native-Cast

Uint

begin

lift-definition *uint-of-uint8* :: *uint8* ⇒ *uint* **is** *ucast* ⟨proof⟩

lift-definition *uint-of-uint16* :: *uint16* ⇒ *uint* **is** *ucast* ⟨proof⟩

lift-definition *uint-of-uint32* :: *uint32* ⇒ *uint* **is** *ucast* ⟨proof⟩

lift-definition *uint-of-uint64* :: *uint64* ⇒ *uint* **is** *ucast* ⟨proof⟩

lift-definition *uint8-of-uint* :: *uint* ⇒ *uint8* **is** *ucast* ⟨proof⟩

lift-definition *uint16-of-uint* :: *uint* ⇒ *uint16* **is** *ucast* ⟨proof⟩

lift-definition *uint32-of-uint* :: *uint* ⇒ *uint32* **is** *ucast* ⟨proof⟩

lift-definition *uint64-of-uint* :: *uint* ⇒ *uint64* **is** *ucast* ⟨proof⟩

code-printing

constant *uint-of-uint8* →

(*SML*) *Word.fromLarge* (*Word8.toLarge* -) **and**

```

(Haskell) (Prelude.fromIntegral - :: Uint.Word) and
(Scala) ((-).toInt & 0xFF)
| constant uint-of-uint16 →
(SML-word) Word.fromLarge (Word16.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint.Word) and
(Scala) (-).toInt
| constant uint-of-uint32 →
(SML) Word.fromLarge (Word32.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint.Word) and
(Scala) - and
(OCaml) (Int32.to'-int -) land Uint.int'-mask
| constant uint-of-uint64 →
(SML) Word.fromLarge (Uint64.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint.Word) and
(Scala) (-).toInt and
(OCaml) Int64.to'-int
| constant uint8-of-uint →
(SML) Word8.fromLarge (Word.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint8.Word8) and
(Scala) (-).toByte
| constant uint16-of-uint →
(SML-word) Word16.fromLarge (Word.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint16.Word16) and
(Scala) (-).toChar
| constant uint32-of-uint →
(SML) Word32.fromLarge (Word.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
(Scala) - and
(OCaml) Int32.logand (Int32.of'-int -) Uint.int32'-mask
| constant uint64-of-uint →
(SML) Uint64.fromLarge (Word.toLarge -) and
(Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
(Scala) ((-).toLong & 0xFFFFFFFFL) and
(OCaml) Int64.logand (Int64.of'-int -) Uint.int64'-mask

```

lemma *uint8-of-uint-code* [code]:
uint8-of-uint $x = \text{Abs-uint8}' (\text{ucast} (\text{Rep-uint}' x))$
 ⟨proof⟩

lemma *uint16-of-uint-code* [code]:
uint16-of-uint $x = \text{Abs-uint16}' (\text{ucast} (\text{Rep-uint}' x))$
 ⟨proof⟩

lemma *uint32-of-uint-code* [code]:
uint32-of-uint $x = \text{Abs-uint32}' (\text{ucast} (\text{Rep-uint}' x))$
 ⟨proof⟩

lemma *uint64-of-uint-code* [code]:
uint64-of-uint $x = \text{Abs-uint64}' (\text{ucast} (\text{Rep-uint}' x))$

<proof>

lemma *wint-of-wint8-code* [code]:

wint-of-wint8 $x = \text{Abs-wint}' (\text{ucast} (\text{Rep-wint8}' x))$

<proof>

lemma *wint-of-wint16-code* [code]:

wint-of-wint16 $x = \text{Abs-wint}' (\text{ucast} (\text{Rep-wint16}' x))$

<proof>

lemma *wint-of-wint32-code* [code]:

wint-of-wint32 $x = \text{Abs-wint}' (\text{ucast} (\text{Rep-wint32}' x))$

<proof>

lemma *wint-of-wint64-code* [code]:

wint-of-wint64 $x = \text{Abs-wint}' (\text{ucast} (\text{Rep-wint64}' x))$

<proof>

end

9.2 Compatibility with Imperative/HOL

theory *Native-Word-Imperative-HOL* **imports**

Code-Target-Word

HOL-Imperative-HOL.Heap-Monad

begin

We add a code target that combines the translations for native words that are by default not supported by all PolyML versions with the adaptations for `Imperative_HOL`.

<ML>

end

Chapter 10

Test cases

```
theory Native-Word-Test
imports
  Uint64 Uint32 Uint16 Uint8 Uint Native-Cast-Uint
  HOL-Library.Code-Test
begin

export-code
  nat-of-uint8 uint8-of-nat
  nat-of-uint16 uint16-of-nat
  nat-of-uint32 uint32-of-nat
  nat-of-uint64 uint64-of-nat
  nat-of-uint uint-of-nat
in SML
```

10.1 Tests for $\text{isa}^{\wedge}\text{typinteger}$

```
context
  includes bit-operations-syntax
begin

definition bit-integer-test :: bool
  where  $\langle \text{bit-integer-test} =$ 
    (([ -1 AND 3, 1 AND -3, 3 AND 5, -3 AND (- 5)
      , -3 OR 1, 1 OR -3, 3 OR 5, -3 OR (- 5)
      , NOT 1, NOT (- 3)
      , -1 XOR 3, 1 XOR (- 3), 3 XOR 5, -5 XOR (- 3)
      , Bit-Operations.set-bit 4 5, Bit-Operations.set-bit 2 (- 5)
      , Bit-Operations.unset-bit 0 5, Bit-Operations.unset-bit 1 (- 5)
      , Bit-Operations.flip-bit 4 5, Bit-Operations.flip-bit 1 (- 5)
      , push-bit 2 1, push-bit 3 (- 1)
      , drop-bit 3 100, drop-bit 3 (- 100)
      , take-bit 4 100, take-bit 4 (- 100)] :: integer list)
```

```

= [ 3, 1, 1, -7
  , -3, -3, 7, -1
  , -2, 2
  , -4, -4, 6, 6
  , 21, -1, 4, -7, 21, -7
  , 4, -8
  , 12, -13
  , 4, 12] ^
[ bit (5 :: integer) 4, bit (5 :: integer) 2, bit (-5 :: integer) 4, bit (-5 ::
integer) 2
, bit (5 :: integer) 0, bit (4 :: integer) 0, bit (-1 :: integer) 0, bit (-2 ::
integer) 0,
  msb (5 :: integer), msb (0 :: integer), msb (-1 :: integer), msb (-2 :: integer)]
= [ False, True, True, False,
  True, False, True, False,
  False, False, True, True]

```

```

export-code bit-integer-test
  checking SML Haskell? Haskell-Quickcheck? OCaml? Scala

```

```

notepad

```

```

begin
  <proof>
end

```

```

<ML>

```

```

lemma <x AND y = x OR (y :: integer)>

```

```

quickcheck [random, expect=counterexample]

```

```

quickcheck [exhaustive, expect=counterexample]

```

```

<proof>

```

```

lemma <(x :: integer) AND x = x OR x>

```

```

quickcheck [narrowing, expect=no-counterexample]

```

```

<proof>

```

```

lemma <(f :: integer => unit) = g>

```

```

quickcheck[narrowing, size=3, expect=no-counterexample]

```

```

<proof>

```

```

end

```

10.2 Tests for *uint8*

```

context

```

```

  includes bit-operations-syntax

```

```

begin

```

```

definition test-uint8 :: bool

```

```

where <test-uint8 <→>
(( [ 0x101, -1, -255, 0xFF, 0x12
  , 0x5A AND 0x36
  , 0x5A OR 0x36
  , 0x5A XOR 0x36
  , NOT 0x5A
  , 5 + 6, -5 + 6, -6 + 5, -5 + -6, 0xFF + 1
  , 5 - 3, 3 - 5
  , 5 * 3, -5 * 3, -5 * -4, 0x12 * 0x87
  , 5 div 3, -5 div 3, -5 div -3, 5 div -3
  , 5 mod 3, -5 mod 3, -5 mod -3, 5 mod -3
  , Bit-Operations.set-bit 4 5, Bit-Operations.set-bit 2 (- 5)
  , Bit-Operations.set-bit 32 5, Bit-Operations.set-bit 32 (- 5)
  , Bit-Operations.unset-bit 0 5, Bit-Operations.unset-bit 1 (- 5)
  , Bit-Operations.unset-bit 32 5, Bit-Operations.unset-bit 32 (- 5)
  , Bit-Operations.flip-bit 4 5, Bit-Operations.flip-bit 1 (- 5)
  , Bit-Operations.flip-bit 32 5, Bit-Operations.flip-bit 32 (- 5)
  , push-bit 2 1, push-bit 3 (- 1), push-bit 8 1, push-bit 0 1
  , drop-bit 3 100, drop-bit 3 (- 100), drop-bit 8 100, drop-bit 8 (- 100)
  , signed-drop-bit-uint8 3 100, signed-drop-bit-uint8 3 (- 100)
  , signed-drop-bit-uint8 8 100, signed-drop-bit-uint8 8 (- 100)
  , take-bit 4 100, take-bit 4 (- 100)] :: uint8 list)
=
[ 1, 255, 1, 255, 18
  , 18
  , 126
  , 108
  , 165
  , 11, 1, 255, 245, 0
  , 2, 254
  , 15, 241, 20, 126
  , 1, 83, 0, 0
  , 2, 2, 251, 5
  , 21, 255, 5, 251, 4, 249, 5, 251, 21, 249, 5, 251
  , 4, 248, 0, 1
  , 12, 19, 0, 0
  , 12, 243, 0, 255
  , 4, 12] ) ^
(( (0x5 :: uint8) = 0x5, (0x5 :: uint8) = 0x6
  , (0x5 :: uint8) < 0x5, (0x5 :: uint8) < 0x6, (-5 :: uint8) < 6, (6 :: uint8) <
-5
  , (0x5 :: uint8) ≤ 0x5, (0x5 :: uint8) ≤ 0x4, (-5 :: uint8) ≤ 6, (6 :: uint8) ≤
-5
  , (0x7F :: uint8) < 0x80, (0xFF :: uint8) < 0, (0x80 :: uint8) < 0x7F
  , bit (0x7F :: uint8) 0, bit (0x7F :: uint8) 7, bit (0x80 :: uint8) 7, bit (0x80 ::
uint8) 8
  ]
=
[ True, False

```

```

, False, True, False, True
, True, False, False, True
, True, False, False
, True, False, True, False
]) ^
([integer-of-uint8 0, integer-of-uint8 0x7F, integer-of-uint8 0x80, integer-of-uint8
0xAA]
=
[0, 0x7F, 0x80, 0xAA])

```

```

export-code test-uint8
  checking SML Haskell? Scala

```

```

notepad
begin
  ⟨proof⟩
end

```

⟨ML⟩

```

definition test-uint8' :: uint8
  where ⟨test-uint8' = drop-bit 2 (push-bit 3 (0 + 10 - 14 * 3 div 6 mod 3))⟩

```

⟨ML⟩

```

lemma ⟨x AND y = x OR (y :: uint8)⟩
quickcheck [random, expect=counterexample]
quickcheck [exhaustive, expect=counterexample]
⟨proof⟩

```

```

lemma ⟨(x :: uint8) AND x = x OR x⟩
quickcheck [narrowing, expect=no-counterexample]
⟨proof⟩

```

```

lemma ⟨(f :: uint8 ⇒ unit) = g⟩
quickcheck [narrowing, size=3, expect=no-counterexample]
⟨proof⟩

```

10.3 Tests for *uint16*

```

context
  includes bit-operations-syntax
begin

```

```

definition test-uint16 :: bool
  where ⟨test-uint16 ↔
    (([ 0x10001, -1, -65535, 0xFFFF, 0x1234
      , 0x5A AND 0x36
      , 0x5A OR 0x36

```

```

, 0x5A XOR 0x36
, NOT 0x5A
, 5 + 6, -5 + 6, -6 + 5, -5 + -6, 0xFFFF + 1
, 5 - 3, 3 - 5
, 5 * 3, -5 * 3, -5 * -4, 0x1234 * 0x8765
, 5 div 3, -5 div 3, -5 div -3, 5 div -3
, 5 mod 3, -5 mod 3, -5 mod -3, 5 mod -3
, Bit-Operations.set-bit 4 5, Bit-Operations.set-bit 2 (- 5)
, Bit-Operations.set-bit 32 5, Bit-Operations.set-bit 32 (- 5)
, Bit-Operations.unset-bit 0 5, Bit-Operations.unset-bit 1 (- 5)
, Bit-Operations.unset-bit 32 5, Bit-Operations.unset-bit 32 (- 5)
, Bit-Operations.flip-bit 4 5, Bit-Operations.flip-bit 1 (- 5)
, Bit-Operations.flip-bit 32 5, Bit-Operations.flip-bit 32 (- 5)
, push-bit 2 1, push-bit 3 (- 1), push-bit 16 1, push-bit 0 1
, drop-bit 3 100, drop-bit 3 (- 100), drop-bit 16 100, drop-bit 16 (- 100)
, signed-drop-bit-uint16 3 100, signed-drop-bit-uint16 3 (- 100)
, signed-drop-bit-uint16 16 100, signed-drop-bit-uint16 16 (- 100)
, take-bit 4 100, take-bit 4 (- 100)] :: uint16 list)
=
[ 1, 65535, 1, 65535, 4660
, 18
, 126
, 108
, 65445
, 11, 1, 65535, 65525, 0
, 2, 65534
, 15, 65521, 20, 39556
, 1, 21843, 0, 0
, 2, 2, 65531, 5
, 21, 65535, 5, 65531, 4, 65529, 5, 65531, 21, 65529, 5, 65531
, 4, 65528, 0, 1
, 12, 8179, 0, 0
, 12, 65523, 0, 65535
, 4, 12]) ^
([ (0x5 :: uint16) = 0x5, (0x5 :: uint16) = 0x6
, (0x5 :: uint16) < 0x5, (0x5 :: uint16) < 0x6, (-5 :: uint16) < 6, (6 :: uint16)
< -5
, (0x5 :: uint16) ≤ 0x5, (0x5 :: uint16) ≤ 0x4, (-5 :: uint16) ≤ 6, (6 :: uint16)
≤ -5
, (0x7FFF :: uint16) < 0x8000, (0xFFFF :: uint16) < 0, (0x8000 :: uint16) <
0x7FFF
, bit (0x7FFF :: uint16) 0, bit (0x7FFF :: uint16) 15, bit (0x8000 :: uint16)
15, bit (0x8000 :: uint16) 16
]
=
[ True, False
, False, True, False, True
, True, False, False, True
, True, False, False

```

```

    , True, False, True, False
  ]) ^
  ([integer-of-uint16 0, integer-of-uint16 0x7FFF, integer-of-uint16 0x8000, inte-
  ger-of-uint16 0xAAAA]
  =
  [0, 0x7FFF, 0x8000, 0xAAAA])

```

```

export-code test-uint16 checking Haskell? Scala
export-code test-uint16 checking SML-word

```

```

notepad begin

```

```

  <proof>

```

```

end

```

```

lemma <(x :: uint16) AND x = x OR x>

```

```

quickcheck [narrowing, expect=no-counterexample]

```

```

  <proof>

```

```

lemma <(f :: uint16 => unit) = g>

```

```

quickcheck [narrowing, size=3, expect=no-counterexample]

```

```

  <proof>

```

```

end

```

10.4 Tests for *uint32*

```

context

```

```

  includes bit-operations-syntax

```

```

begin

```

```

definition test-uint32 :: bool

```

```

  where <test-uint32 <=>>

```

```

  (([ 0x100000001, -1, -4294967291, 0xFFFFFFFF, 0x12345678
    , 0x5A AND 0x36
    , 0x5A OR 0x36
    , 0x5A XOR 0x36
    , NOT 0x5A
    , 5 + 6, -5 + 6, -6 + 5, -5 + (- 6), 0xFFFFFFFF + 1
    , 5 - 3, 3 - 5
    , 5 * 3, -5 * 3, -5 * -4, 0x12345678 * 0x87654321
    , 5 div 3, -5 div 3, -5 div -3, 5 div -3
    , 5 mod 3, -5 mod 3, -5 mod -3, 5 mod -3
    , Bit-Operations.set-bit 4 5, Bit-Operations.set-bit 2 (- 5)
    , Bit-Operations.set-bit 32 5, Bit-Operations.set-bit 32 (- 5)
    , Bit-Operations.unset-bit 0 5, Bit-Operations.unset-bit 1 (- 5)
    , Bit-Operations.unset-bit 32 5, Bit-Operations.unset-bit 32 (- 5)
    , Bit-Operations.flip-bit 4 5, Bit-Operations.flip-bit 1 (- 5)
    , Bit-Operations.flip-bit 32 5, Bit-Operations.flip-bit 32 (- 5)
    , push-bit 2 1, push-bit 3 (- 1), push-bit 32 1, push-bit 0 1

```

```

, drop-bit 3 100, drop-bit 3 (- 100), drop-bit 32 100, drop-bit 32 (- 100)
, signed-drop-bit-uint32 3 100, signed-drop-bit-uint32 3 (- 100)
, signed-drop-bit-uint32 32 100, signed-drop-bit-uint32 32 (- 100)
, take-bit 4 100, take-bit 4 (- 100)] :: uint32 list)
=
[ 1, 4294967295, 5, 4294967295, 305419896
, 18
, 126
, 108
, 4294967205
, 11, 1, 4294967295, 4294967285, 0
, 2, 4294967294
, 15, 4294967281, 20, 1891143032
, 1, 1431655763, 0, 0
, 2, 2, 4294967291, 5
, 21, 4294967295, 5, 4294967291, 4, 4294967289, 5, 4294967291, 21, 4294967289,
5, 4294967291
, 4, 4294967288, 0, 1
, 12, 536870899, 0, 0
, 12, 4294967283, 0, 4294967295
, 4, 12]) ^
([ (0x5 :: uint32) = 0x5, (0x5 :: uint32) = 0x6
, (0x5 :: uint32) < 0x5, (0x5 :: uint32) < 0x6, (-5 :: uint32) < 6, (6 :: uint32)
< -5
, (0x5 :: uint32) ≤ 0x5, (0x5 :: uint32) ≤ 0x4, (-5 :: uint32) ≤ 6, (6 :: uint32)
≤ -5
, (0x7FFFFFFF :: uint32) < 0x80000000, (0xFFFFFFFF :: uint32) < 0,
(0x80000000 :: uint32) < 0x7FFFFFFF
, bit (0x7FFFFFFF :: uint32) 0, bit (0x7FFFFFFF :: uint32) 31, bit (0x80000000
:: uint32) 31, bit (0x80000000 :: uint32) 32
]
)
=
[ True, False
, False, True, False, True
, True, False, False, True
, True, False, False
, True, False, True, False
] ^
([integer-of-uint32 0, integer-of-uint32 0x7FFFFFFF, integer-of-uint32 0x80000000,
integer-of-uint32 0xAFFFFFFF])
=
[0, 0x7FFFFFFF, 0x80000000, 0xAFFFFFFF]

```

export-code *test-uint32 checking SML Haskell? OCaml? Scala*

notepad begin

<proof>

end

⟨ML⟩

definition *test-wint32'* :: *wint32*
where $\langle \text{test-wint32}' = \text{drop-bit } 2 \ (\text{push-bit } 3 \ (0 + 10 - 14 * 3 \ \text{div } 6 \ \text{mod } 3)) \rangle$

⟨ML⟩

lemma $x \ \text{AND} \ y = x \ \text{OR} \ (y :: \text{wint32})$
quickcheck [*random, expect=counterexample*]
quickcheck [*exhaustive, expect=counterexample*]
 ⟨*proof*⟩

lemma $(x :: \text{wint32}) \ \text{AND} \ x = x \ \text{OR} \ x$
quickcheck [*narrowing, expect=no-counterexample*]
 ⟨*proof*⟩

lemma $(f :: \text{wint32} \Rightarrow \text{unit}) = g$
quickcheck [*narrowing, size=3, expect=no-counterexample*]
 ⟨*proof*⟩

end

10.5 Tests for *wint64*

context
includes *bit-operations-syntax*
begin

definition *test-wint64* :: *bool*
where $\langle \text{test-wint64} \longleftrightarrow$
 (([*0x10000000000000001*, -1 , -9223372036854775808 , *0xFFFFFFFFFFFFFFFF*,
0x1234567890ABCDEF
 , $0x5A \ \text{AND} \ 0x36$
 , $0x5A \ \text{OR} \ 0x36$
 , $0x5A \ \text{XOR} \ 0x36$
 , $\text{NOT } 0x5A$
 , $5 + 6$, $-5 + 6$, $-6 + 5$, $-5 + (-6)$, $0xFFFFFFFFFFFFFFFF + 1$
 , $5 - 3$, $3 - 5$
 , $5 * 3$, $-5 * 3$, $-5 * -4$, $0x1234567890ABCDEF * 0xFEDCBA0987654321$
 , $5 \ \text{div} \ 3$, $-5 \ \text{div} \ 3$, $-5 \ \text{div} \ -3$, $5 \ \text{div} \ -3$
 , $5 \ \text{mod} \ 3$, $-5 \ \text{mod} \ 3$, $-5 \ \text{mod} \ -3$, $5 \ \text{mod} \ -3$
 , *Bit-Operations.set-bit* 4 5, *Bit-Operations.set-bit* 2 (- 5)
 , *Bit-Operations.set-bit* 32 5, *Bit-Operations.set-bit* 32 (- 5)
 , *Bit-Operations.unset-bit* 0 5, *Bit-Operations.unset-bit* 1 (- 5)
 , *Bit-Operations.unset-bit* 32 5, *Bit-Operations.unset-bit* 32 (- 5)
 , *Bit-Operations.flip-bit* 4 5, *Bit-Operations.flip-bit* 1 (- 5)
 , *Bit-Operations.flip-bit* 32 5, *Bit-Operations.flip-bit* 32 (- 5)
 , *push-bit* 2 1, *push-bit* 3 (- 1), *push-bit* 64 1, *push-bit* 0 1
 , *drop-bit* 3 100, *drop-bit* 3 (- 100), *drop-bit* 64 100, *drop-bit* 64 (- 100)

```

, signed-drop-bit-uint64 3 100, signed-drop-bit-uint64 3 (- 100)
, signed-drop-bit-uint64 64 100, signed-drop-bit-uint64 64 (- 100)
, take-bit 4 100, take-bit 4 (- 100)] :: uint64 list)
=
[ 1, 18446744073709551615, 9223372036854775808, 18446744073709551615,
1311768467294899695
, 18
, 126
, 108
, 18446744073709551525
, 11, 1, 18446744073709551615, 18446744073709551605, 0
, 2, 18446744073709551614
, 15, 18446744073709551601, 20, 14000077364136384719
, 1, 6148914691236517203, 0, 0
, 2, 2, 18446744073709551611, 5
, 21, 18446744073709551615, 4294967301, 18446744073709551611, 4, 18446744073709551609,
5, 18446744069414584315, 21, 18446744073709551609, 4294967301, 18446744069414584315
, 4, 18446744073709551608, 0, 1
, 12, 2305843009213693939, 0, 0
, 12, 18446744073709551603, 0, 18446744073709551615
, 4, 12]) ^
([ (0x5 :: uint64) = 0x5, (0x5 :: uint64) = 0x6
, (0x5 :: uint64) < 0x5, (0x5 :: uint64) < 0x6, (-5 :: uint64) < 6, (6 :: uint64)
< -5
, (0x5 :: uint64) ≤ 0x5, (0x5 :: uint64) ≤ 0x4, (-5 :: uint64) ≤ 6, (6 :: uint64)
≤ -5
, (0x7FFFFFFFFFFFFFFFFF :: uint64) < 0x8000000000000000, (0xFFFFFFFFFFFFFFFF
:: uint64) < 0, (0x8000000000000000 :: uint64) < 0x7FFFFFFFFFFFFFFFFF
, bit (0x7FFFFFFFFFFFFFFFFF :: uint64) 0, bit (0x7FFFFFFFFFFFFFFFFF ::
uint64) 63, bit (0x8000000000000000 :: uint64) 63, bit (0x8000000000000000 ::
uint64) 64
]
=
[ True, False
, False, True, False, True
, True, False, False, True
, True, False, False
, True, False, True, False
]) ^
([integer-of-uint64 0, integer-of-uint64 0x7FFFFFFFFFFFFFFFFF, integer-of-uint64
0x8000000000000000, integer-of-uint64 0xAAAAAAAAAAAAAAAAAA])
=
[0, 0x7FFFFFFFFFFFFFFFFF, 0x8000000000000000, 0xAAAAAAAAAAAAAAAAAA])
value [nbe] <[0x10000000000000001, -1, -9223372036854775808, 0xFFFFFFFFFFFFFFFF,
0x1234567890ABCDEF
, 0x5A AND 0x36
, 0x5A OR 0x36
, 0x5A XOR 0x36

```

```

, NOT 0x5A
, 5 + 6, -5 + 6, -6 + 5, -5 + (- 6), 0xFFFFFFFFFFFFFFFF + 1
, 5 - 3, 3 - 5
, 5 * 3, -5 * 3, -5 * -4, 0x1234567890ABCDEF * 0xFEDCBA0987654321
, 5 div 3, -5 div 3, -5 div -3, 5 div -3
, 5 mod 3, -5 mod 3, -5 mod -3, 5 mod -3
, push-bit 2 1, push-bit 3 (- 1), push-bit 64 1, push-bit 0 1
, drop-bit 3 100, drop-bit 3 (- 100), drop-bit 64 100, drop-bit 64 (- 100)
, signed-drop-bit-uint64 3 100, signed-drop-bit-uint64 3 (- 100)
, signed-drop-bit-uint64 64 100, signed-drop-bit-uint64 64 (- 100)
, take-bit 4 100, take-bit 4 (- 100)] :: uint64 list

```

export-code *test-uint64* **checking** *SML Haskell? OCaml? Scala*

notepad begin

⟨*proof*⟩

end

⟨*ML*⟩

definition *test-uint64'* :: *uint64*

where ⟨*test-uint64' = drop-bit 2 (push-bit 3 (0 + 10 - 14 * 3 div 6 mod 3))*⟩

⟨*ML*⟩

end

10.6 Tests for *uint*

context

includes *bit-operations-syntax*

begin

definition *test-uint* :: *bool*

where ⟨*test-uint = (let*

test-list1 = (let

HS = uint-of-int (2 ^ (dflt-size - 1))

in

[*HS + HS + 1, -1, -HS - HS + 5, HS + (HS - 1), 0x12*

, 0x5A AND 0x36

, 0x5A OR 0x36

, 0x5A XOR 0x36

, NOT 0x5A

, 5 + 6, -5 + 6, -6 + 5, -5 + -6, HS + (HS - 1) + 1

, 5 - 3, 3 - 5

*, 5 * 3, -5 * 3, -5 * -4, 0x12345678 * 0x87654321]*

Ⓢ (if *dflt-size > 4* then

[*5 div 3, -5 div 3, -5 div -3, 5 div -3*

, 5 mod 3, -5 mod 3, -5 mod -3, 5 mod -3

```

, Bit-Operations.set-bit dflt-size 5, Bit-Operations.set-bit dflt-size (- 5)
, Bit-Operations.unset-bit dflt-size 5, Bit-Operations.unset-bit dflt-size (- 5)
, Bit-Operations.flip-bit 0 5, Bit-Operations.flip-bit 0 (- 5)
, push-bit 2 1, push-bit 3 (- 1), push-bit dflt-size 1, push-bit 0 1
, drop-bit 3 31, drop-bit 3 (- 1), drop-bit dflt-size 31, drop-bit dflt-size (- 1)
, signed-drop-bit-uint 2 15, signed-drop-bit-uint 3 (- 1)
, signed-drop-bit-uint dflt-size 15, signed-drop-bit-uint dflt-size (- 1)
, take-bit 4 100, take-bit 4 (- 100)]
else [] :: uint list));

```

```

test-list2 = (let
  S = wivs-shift
in
  ([ 1, -1, -S + 5, S - 1, 0x12
    , 0x5A AND 0x36
    , 0x5A OR 0x36
    , 0x5A XOR 0x36
    , NOT 0x5A
    , 5 + 6, -5 + 6, -6 + 5, -5 + -6, 0
    , 5 - 3, 3 - 5
    , 5 * 3, -5 * 3, -5 * -4, 0x12345678 * 0x87654321]
  @ (if dflt-size > 4 then
    [ 5 div 3, (S - 5) div 3, (S - 5) div (S - 3), 5 div (S - 3)
    , 5 mod 3, (S - 5) mod 3, (S - 5) mod (S - 3), 5 mod (S - 3)
    , 5, -5, 5, -5, 4, -6
    , 4, -8, 0, 1
    , 3, drop-bit 3 S - 1, 0, 0
    , 3, drop-bit 1 S + drop-bit 1 S - 1, 0, -1
    , 4, 12]
  else [] :: int list));

```

```

test-list-c1 = (let
  HS = uint-of-int ((2^(dflt-size - 1)))
in
  [ (0x5 :: uint) = 0x5, (0x5 :: uint) = 0x6
    , (0x5 :: uint) < 0x5, (0x5 :: uint) < 0x6, (-5 :: uint) < 6, (6 :: uint) < -5
    , (0x5 :: uint) ≤ 0x5, (0x5 :: uint) ≤ 0x4, (-5 :: uint) ≤ 6, (6 :: uint) ≤ -5
    , (HS - 1) < HS, (HS + HS - 1) < 0, HS < HS - 1
    , bit (HS - 1) 0, bit (HS - 1 :: uint) (dflt-size - 1), bit (HS :: uint) (dflt-size
- 1), bit (HS :: uint) dflt-size
  ]);

```

```

test-list-c2 =
[ True, False
, False, dflt-size ≥ 2, dflt-size = 3, dflt-size ≠ 3
, True, False, dflt-size = 3, dflt-size ≠ 3
, True, False, False
, dflt-size ≠ 1, False, True, False
]

```

```

in
  test-list1 = map uint-of-int test-list2
  ∧ test-list-c1 = test-list-c2)

```

export-code *test-uint checking SML Haskell? OCaml? Scala*

```

lemma test-uint
quickcheck [exhaustive, expect=no-counterexample]
  ⟨proof⟩

```

```

lemma ⟨x AND y = x OR (y :: uint)⟩
quickcheck [random, expect=counterexample]
quickcheck [exhaustive, expect=counterexample]
  ⟨proof⟩

```

```

lemma ⟨(x :: uint) AND x = x OR x⟩
quickcheck [narrowing, expect=no-counterexample]
  ⟨proof⟩

```

```

lemma ⟨(f :: uint ⇒ unit) = g⟩
quickcheck [narrowing, size=3, expect=no-counterexample]
  ⟨proof⟩

```

10.7 Tests for casts

```

definition test-casts :: bool
  where ⟨test-casts ↔
    map uint8-of-uint32 [10, 0, 0xFE, 0xFFFFFFFF] = [10, 0, 0xFE, 0xFF] ∧
    map uint8-of-uint64 [10, 0, 0xFE, 0xFFFFFFFFFFFFFFFF] = [10, 0, 0xFE,
0xFF] ∧
    map uint32-of-uint8 [10, 0, 0xFF] = [10, 0, 0xFF] ∧
    map uint64-of-uint8 [10, 0, 0xFF] = [10, 0, 0xFF]⟩

```

```

definition test-casts' :: bool
  where ⟨test-casts' ↔
    map uint8-of-uint16 [10, 0, 0xFE, 0xFFFF] = [10, 0, 0xFE, 0xFF] ∧
    map uint16-of-uint8 [10, 0, 0xFF] = [10, 0, 0xFF] ∧
    map uint16-of-uint32 [10, 0, 0xFFFE, 0xFFFFFFFF] = [10, 0, 0xFFFE, 0xFFFF]
  ∧
    map uint16-of-uint64 [10, 0, 0xFFFE, 0xFFFFFFFFFFFFFFFF] = [10, 0,
0xFFFE, 0xFFFF] ∧
    map uint32-of-uint16 [10, 0, 0xFFFF] = [10, 0, 0xFFFF] ∧
    map uint64-of-uint16 [10, 0, 0xFFFF] = [10, 0, 0xFFFF]⟩

```

```

definition test-casts'' :: bool
  where ⟨test-casts'' ↔
    map uint32-of-uint64 [10, 0, 0xFFFFFFFFFE, 0xFFFFFFFFFFFFFFFF] = [10,
0, 0xFFFFFFFFFE, 0xFFFFFFFF] ∧
    map uint64-of-uint32 [10, 0, 0xFFFFFFFF] = [10, 0, 0xFFFFFFFF]⟩

```

```

export-code test-casts test-casts'' checking SML Haskell? Scala
export-code test-casts'' checking OCaml?
export-code test-casts' checking Haskell? Scala

```

```

notepad begin

```

```

  ⟨proof⟩

```

```

end

```

```

⟨ML⟩

```

```

definition test-casts-uint :: bool

```

```

  where ⟨test-casts-uint ↔
    map uint-of-uint32 ([0, 10] @ (if dflt-size < 32 then [push-bit (dflt-size - 1) 1,
0xFFFFFFFF] else [0xFFFFFFFF])) =
    [0, 10] @ (if dflt-size < 32 then [push-bit (dflt-size - 1) 1, (push-bit dflt-size 1)
- 1] else [0xFFFFFFFF]) ∧
    map uint32-of-uint [0, 10, if dflt-size < 32 then push-bit (dflt-size - 1) 1 else
0xFFFFFFFF] =
    [0, 10, if dflt-size < 32 then push-bit (dflt-size - 1) 1 else 0xFFFFFFFF] ∧
    map uint-of-uint64 [0, 10, push-bit (dflt-size - 1) 1, 0xFFFFFFFFFFFFFFFF]
  =
    [0, 10, push-bit (dflt-size - 1) 1, (push-bit dflt-size 1) - 1] ∧
    map uint64-of-uint [0, 10, push-bit (dflt-size - 1) 1] =
    [0, 10, push-bit (dflt-size - 1) 1]⟩

```

```

definition test-casts-uint' :: bool

```

```

  where ⟨test-casts-uint' ↔
    map uint-of-uint16 [0, 10, 0xFFFF] = [0, 10, 0xFFFF] ∧
    map uint16-of-uint [0, 10, 0xFFFF] = [0, 10, 0xFFFF]⟩

```

```

definition test-casts-uint'' :: bool

```

```

  where ⟨test-casts-uint'' ↔
    map uint-of-uint8 [0, 10, 0xFF] = [0, 10, 0xFF] ∧
    map uint8-of-uint [0, 10, 0xFF] = [0, 10, 0xFF]⟩

```

```

export-code test-casts-uint test-casts-uint'' checking SML Haskell?

```

```

export-code test-casts-uint checking OCaml?

```

```

export-code test-casts-uint' checking Haskell? Scala

```

```

end

```

```

end

```

```

end

```


Chapter 11

Implementation of bit operations on int by target language operations

```
theory Code-Target-Int-Bit
  imports
    HOL-Library.Code-Target-Int
    HOL-Library.Code-Target-Bit-Shifts
    Code-Int-Integer-Conversion
begin

lemma int-of-integer-symbolic-code [code drop: int-of-integer-symbolic, code]:
  int-of-integer-symbolic = int-of-integer
  ⟨proof⟩

context
  includes bit-operations-syntax
begin

context
begin

qualified definition even :: ⟨int ⇒ bool⟩
  where [code-abbrev]: ⟨even = Parity.even⟩

end

lemma [code]:
  ⟨Code-Target-Int-Bit.even i ⟷ i AND 1 = 0⟩
  ⟨proof⟩

lemma [code-unfold]:
  ⟨of-bool (odd i) = i AND 1⟩ for i :: int
```

<proof>

lemma [*code-unfold*]:
<bit x n \longleftrightarrow x AND (push-bit n 1) \neq 0> **for** *x :: int*
<proof>

end

end

theory *Native-Word-Test-Emu*
imports
Native-Word-Test
Code-Target-Int-Bit
begin

11.1 Test cases for emulation of native words

11.1.1 Tests for *wint8*

Test that *wint8* is emulated for OCaml via *8 word* if *Native-Word.Code-Target-Int-Bit* is imported.

definition *test-wint8-emulation* :: *bool*
where *<test-wint8-emulation \longleftrightarrow (0xFFF - 0x10 = (0xEF :: wint8))>*

export-code *test-wint8-emulation* **checking** *OCaml?*
 — test the other target languages as well *SML Haskell? Scala*

11.1.2 Tests for *wint16*

Test that *wint16* is emulated for PolyML and OCaml via *16 word* if *Native-Word.Code-Target-Int-Bit* is imported.

definition *test-wint16-emulation* :: *bool*
where *<test-wint16-emulation \longleftrightarrow (0xFFFFF - 0x1000 = (0xEFFF :: wint16))>*

export-code *test-wint16-emulation* **checking** *SML OCaml?*
 — test the other target languages as well *Haskell? Scala*

notepad **begin**
<proof>
end

<ML>

lemma *<x AND y = x OR (y :: wint16)>*
quickcheck [*random, expect=counterexample*]

```
quickcheck [exhaustive, expect=counterexample]
<proof>
```

```
end
```

```
theory Native-Word-Test-PolyML
imports
  Native-Word-Test
begin
```

11.2 Test with PolyML

```
test-code
  test-uint64 <test-uint64' = 0x12>
  test-uint32 <test-uint32' = 0x12>
  test-uint8 <test-uint8' = 0x12>
  test-uint
  test-casts test-casts''
  test-casts-uint test-casts-uint''
in PolyML
```

```
end
```

```
theory Native-Word-Test-PolyML2
imports
  Native-Word-Test-Emu
begin
```

```
test-code
  test-uint16 test-uint16-emulation
  test-casts'
  test-casts-uint'
in PolyML
```

```
end
```

```
theory Native-Word-Test-PolyML64
imports
  Native-Word-Test
begin
```

```
test-code <test-uint64' = 0x12>
in PolyML
```

```
<ML>
```

end

```
theory Native-Word-Test-Scala
imports
  Native-Word-Test
begin
```

11.3 Test with Scala

In Scala, *uint* and *uint32* are both implemented as type `Int`. When they are used in the same generated program, we have to suppress the type class instances for one of them.

code-printing class-instance *uint32* :: *equal* \rightarrow (*Scala*) –

```
test-code
  test-uint64  $\langle$ test-uint64' = 0x12 $\rangle$ 
  test-uint32  $\langle$ test-uint32' = 0x12 $\rangle$ 
  test-uint16
  test-uint8  $\langle$ test-uint8' = 0x12 $\rangle$ 
  test-uint
  test-casts test-casts' test-casts''
  test-casts-uint test-casts-uint' test-casts-uint''
in Scala

end
```

```
theory Native-Word-Test-MLton
imports
  Native-Word-Test
begin
```

11.4 Test with MLton

```
test-code
  test-uint64  $\langle$ test-uint64' = 0x12 $\rangle$ 
  test-uint32  $\langle$ test-uint32' = 0x12 $\rangle$ 
  test-uint8  $\langle$ test-uint8' = 0x12 $\rangle$ 
  test-uint
  test-casts
  test-casts''
  test-casts-uint
  test-casts-uint''
in MLton
```

MLton provides `Word16` and `Word64` structures. To test them in the `SML_word`

target, we have to associate a driver with the combination.

<ML>

test-code

test-uint64 <test-uint64' = 0x12>

test-uint32 <test-uint32' = 0x12>

test-uint16

test-uint8 <test-uint8' = 0x12>

test-uint

test-casts

test-casts'

test-casts''

test-casts-uint

test-casts-uint'

test-casts-uint''

in *MLton-word*

end

theory *Native-Word-Test-MLton2*

imports

Native-Word-Test-Emu

begin

export-code *test-casts'* **in** *SML module-name* *Generated-Code*

test-code

test-uint16 test-uint16-emulation

test-casts'

test-casts-uint'

in *MLton*

end

Chapter 12

User guide for native words

This tutorial explains how to best use the types for native words like *wint32* in your formalisation. You can base your formalisation

1. either directly on these types,
2. or on the generic *'a word* and only introduce native words a posteriori via code generator refinement.

The first option causes the least overhead if you have to prove only little about the words you use and start a fresh formalisation. Just use the native type *wint32* instead of *32 word* and similarly for *wint64*, *wint16*, and *wint8*. As native word types are meant only for code generation, the lemmas about *'a word* have not been duplicated, but you can transfer theorems between native word types and *'a word* using the transfer package.

Note, however, that this option restricts your work a bit: your own functions cannot be “polymorphic” in the word length, but you have to define a separate function for every word length you need.

The second option is recommended if you already have a formalisation based on *'a word* or if your proofs involve words and their properties. It separates code generation from modelling and proving, i.e., you can work with words as usual. Consequently, you have to manually setup the code generator to use the native types wherever you want. The following describes how to achieve this with moderate effort.

Note, however, that some target languages of the code generator (especially OCaml) do not support all the native word types provided. Therefore, you should only import those types that you need – the theory file for each type mentions at the top the restrictions for code generation. For example, PolyML does not provide the *Word16* structure, and OCaml provides neither *Word8* nor *Word16*. You can still use these theories provided that you also import the theory *Native-Word.Code-Target-Int-Bit* (which implements

int by target-language integers), but these words will be implemented via Isabelle's *Word* library, i.e., you do not gain anything in terms of efficiency.

There is a separate code target *SML-word* for *SML*. If you use one of the native words that PolyML does not support (such as *uint16* and *uint64* in 32-bit mode), but would like to map its operations to the Standard Basis Library functions, make sure to use the target *SML-word* instead of *SML*; if you only use native word sizes that PolyML supports, you can stick with *SML*. This ensures that code generation within Isabelle as used by *Quickcheck*, *value* and `@{code}` in ML blocks continues to work.

12.1 Lifting functions from 'a word to native words

This section shows how to convert functions from 'a word to native words. For example, the following function *sum-squares* computes the sum of the first n square numbers in 16 bit arithmetic using a tail-recursive function *gen-sum-squares* with accumulator; for convenience, *sum-squares-int* takes an integer instead of a word.

function *gen-sum-squares* :: 16 word \Rightarrow 16 word \Rightarrow 16 word **where**
gen-sum-squares accum n =
*(if n = 0 then accum else gen-sum-squares (accum + n * n) (n - 1))* <proof> <proof>
definition *sum-squares* :: 16 word \Rightarrow 16 word **where**
sum-squares = gen-sum-squares 0

definition *sum-squares-int* :: int \Rightarrow 16 word **where**
sum-squares-int n = sum-squares (word-of-int n)

The generated code for *sum-squares* and *sum-squares-int* emulates words with unbounded integers and explicit modulus as specified in the theory *HOL-Library.Word*. But for efficiency, we want that the generated code uses machine words and machine arithmetic. Unfortunately, as 'a word is polymorphic in the word length, the code generator can only do this if we use another type for machine words. The theory *Native-Word.Uint16* defines the type *uint16* for machine words of 16 bits. We just have to follow two steps to use it:

First, we lift all our functions from 16 word to *uint16*, i.e., *sum-squares*, *gen-sum-squares*, and *sum-squares-int* in our case. The theory *Native-Word.Uint16* sets up the lifting package for this and has already taken care of the arithmetic and bit-wise operations.

lift-definition *gen-sum-squares-uint* :: uint16 \Rightarrow uint16 \Rightarrow uint16
is *gen-sum-squares* <proof>

lift-definition *sum-squares-uint* :: uint16 \Rightarrow uint16 **is** *sum-squares* <proof>

lift-definition *sum-squares-int-uint* :: int \Rightarrow uint16 **is** *sum-squares-int* <proof>

Second, we also have to transfer the code equations for our functions. The

attribute *Transfer.transferred* takes care of that, but it is better to check that the transfer succeeded: inspect the theorem to check that the new constants are used throughout.

```
lemmas [Transfer.transferred, code] =
  gen-sum-squares.simps
  sum-squares-def
  sum-squares-int-def
```

Finally, we export the code to standard ML. We use the target *SML-word* instead of *SML* to have the operations on *uint16* mapped to the Standard Basis Library. As PolyML does not provide a *Word16* type, the mapping for *uint16* is only active in the refined target *SML-word*.

```
export-code sum-squares-int-uint in SML-word
```

Nevertheless, we can still evaluate terms with *uint16* within Isabelle, i.e., PolyML, but this will be translated to *16 word* and therefore less efficient.

```
value sum-squares-int-uint 40
```

12.2 Storing native words in datatypes

The above lifting is necessary for all functions whose type mentions the word type. Fortunately, we do not have to duplicate functions that merely operate on datatypes that contain words. Nevertheless, we have to tell the code generator that these functions should call the new ones, which operate on machine words. This section shows how to achieve this with data refinement.

12.2.1 Example: expressions and two semantics

As the running example, we consider a language of expressions (literal values, less-than comparisons and conditional) where values are either booleans or 32-bit words. The original specification uses the type *32 word*.

```
datatype val = Bool bool | Word 32 word
datatype expr = Lit val | LT expr expr | IF expr expr expr
```

```
abbreviation (input) word :: 32 word ⇒ expr where word i ≡ Lit (Word i)
```

```
abbreviation (input) bool :: bool ⇒ expr where bool i ≡ Lit (Bool i)
```

— Denotational semantics of expressions, *None* denotes a type error

```
fun eval :: expr ⇒ val option where
  eval (Lit v) = Some v
| eval (LT e1 e2) =
  (case (eval e1, eval e2)
   of (Some (Word i1), Some (Word i2)) ⇒ Some (Bool (i1 < i2))
   | - ⇒ None)
| eval (IF e1 e2 e3) =
```

(*case eval* e_1 of *Some* (*Bool* b) \Rightarrow if b then *eval* e_2 else *eval* e_3
 | - \Rightarrow *None*)

— Small-step semantics of expressions, it gets stuck upon type errors.

inductive step :: *expr* \Rightarrow *expr* \Rightarrow *bool* (- \rightarrow - [50, 50] 60) **where**
 $e \rightarrow e' \Longrightarrow$ *LT* e $e_2 \rightarrow$ *LT* e' e_2
 | $e \rightarrow e' \Longrightarrow$ *LT* (*word* i) $e \rightarrow$ *LT* (*word* i) e'
 | *LT* (*word* i_1) (*word* i_2) \rightarrow *bool* ($i_1 < i_2$)
 | $e \rightarrow e' \Longrightarrow$ *IF* e e_1 $e_2 \rightarrow$ *IF* e' e_1 e_2
 | *IF* (*bool* *True*) e_1 $e_2 \rightarrow$ e_1
 | *IF* (*bool* *False*) e_1 $e_2 \rightarrow$ e_2

— Compile the inductive definition with the predicate compiler

code-pred (*modes*: $i \Rightarrow o \Rightarrow$ *bool* as *reduce*, $i \Rightarrow i \Rightarrow$ *bool* as *step'*) *step* \langle *proof* \rangle

12.2.2 Change the datatype to use machine words

Now, we want to use *uint32* instead of *32 word*. The goal is to make the code generator use the new type without duplicating any of the types (*val*, *expr*) or the functions (*eval*, *reduce*) on such types.

The constructor *Word* has *32 word* in its type, so we have to lift it to *Word'*, and the same holds for the case combinator *case-val*, which *case-val'* replaces.¹ Next, we set up the code generator accordingly: *Bool* and *Word'* are the new constructors for *val*, and *case-val'* is the new case combinator with an appropriate case certificate.² We delete the code equations for the old constructor *Word* and case combinator *case-val* such that the code generator reports missing adaptations.

lift-definition *Word'* :: *uint32* \Rightarrow *val* **is** *Word* \langle *proof* \rangle

code-datatype *Bool* *Word'*

¹Note that we should not declare a case translation for the new case combinator because this will break parsing case expressions with old case combinator.

²Case certificates tell the code generator to replace the HOL case combinator for a datatype with the case combinator of the target language. Without a case certificate, the code generator generates a function that re-implements the case combinator; in a strict languages like ML or Scala, this means that the code evaluates all possible cases before it decides which one is taken.

Case certificates are described in Haftmann's PhD thesis [1, Def. 27]. For a datatype *dt* with constructors C_1 to C_n where each constructor C_i takes k_i parameters, the certificate for the case combinator *case-dt* looks as follows:

lemma

assumes *CASE* \equiv *dt-case* c_1 c_2 ... c_n

shows (*CASE* (C_1 a_{11} a_{12} ... a_{1k_1}) \equiv c_1 a_{11} a_{12} ... a_{1k_1})

&&& (*CASE* (C_2 a_{21} a_{22} ... a_{2k_2}) \equiv c_2 a_{21} a_{22} ... a_{2k_2})

&&& ...

&&& (*CASE* (C_n a_{n1} a_{n2} ... a_{nk_n}) \equiv c_n a_{n1} a_{n2} ... a_{nk_n})

lift-definition *case-val'* :: (bool \Rightarrow 'a) \Rightarrow (uint32 \Rightarrow 'a) \Rightarrow val \Rightarrow 'a **is** *case-val*
 <proof>

lemmas [*code*, *simp*] = val.case [Transfer.transferred]

lemma *case-val'-cert*:

fixes bool word' b w

assumes CASE \equiv *case-val'* bool word'

shows (CASE (Bool b) \equiv bool b) &&& (CASE (Word' w) \equiv word' w)

<proof>

<ML>

declare [[*code drop*: *case-val* Word]]

12.2.3 Make functions use functions on machine words

Finally, we merely have to change the code equations to use the new functions that operate on *uint32*. As before, the attribute *Transfer.transferred* does the job. In our example, we adapt the equality test on *val* (code equations *val.eq.simps*) and the denotational and small-step semantics (code equations *eval.simps* and *step.equation*, respectively).

We check that the adaptation has succeeded by exporting the functions. As we only use native word sizes that PolyML supports, we can use the usual target *SML* instead of *SML-word*.

lemmas [*code*] =

val.eq.simps[*THEN meta-eq-to-obj-eq*, *Transfer.transferred*, *THEN eq-reflection*]

eval.simps[*Transfer.transferred*]

step.equation[*Transfer.transferred*]

export-code *reduce step' eval checking SML*

12.3 Troubleshooting

This section explains some possible problems when using native words. If you experience other difficulties, please contact the author.

12.3.1 *export-code* raises an exception

Probably, you have defined and are using a function on a native word type, but the code equation refers to emulated words. For example, the following defines a function *double* that doubles a word. When we try to export code for *double* without any further setup, *export-code* will raise an exception or generate code that does not compile.

lift-definition *double* :: uint32 \Rightarrow uint32 **is** $\lambda x. x + x$ <proof>

We have to prove a code equation that only uses the existing operations on *uint32*. Then, *export-code* works again.

lemma *double-code* [*code*]: *double n = n + n*
<proof>

12.3.2 The generated code does not compile

Probably, you have been exporting to a target language for which there is no setup, or your compiler does not provide the required API. Every theory for native words mentions at the start the limitations on code generation. Check that your concrete application meets all the requirements.

Alternatively, this might be an instance of the problem described in §12.3.1. For Haskell, you have to enable the extension `TypeSynonymInstances` with `-XTypeSynonymInstances` if you are using polymorphic bit operations on the native word types.

12.3.3 The generated code is too slow

The generated code will most likely not be as fast as a direct implementation in the target language with manual tuning. This is because we want the configuration of the code generation to be sound (as it can be used to prove theorems in Isabelle). Therefore, the bit operations sometimes perform range checks before they call the target language API. Here are some examples:

- Shift distances and bit indices in target languages are often expected to fit into a bounded integer or word. However, the size of these types varies across target languages and platforms. Hence, no Isabelle/HOL type can model uniformly all of them. Instead, the bit operations use arbitrary-precision integers for such quantities and check at run-time that the values fit into a bounded integer or word, respectively – if not, they raise an exception.
- Division and modulo operations explicitly test whether the divisor is 0 and return the HOL value of division by 0 in that case. This is necessary because some languages leave the behaviour of division by 0 unspecified.

If you have better ideas how to eliminate such checks and speed up the generated code without sacrificing soundness, please contact the author!

Bibliography

- [1] F. Haftmann. *Code Generation from Specifications in Higher-Order-Logic*. PhD thesis, Institut für Informatik, Technische Universität München, 2009.