Native Words

Andreas Lochbihler

December 14, 2021

Abstract

This entry makes machine words and machine arithmetic available for code generation from Isabelle/HOL. It provides a common abstraction that hides the differences between the different target languages. The code generator maps these operations to the APIs of the target languages. Apart from that, we extend the available bit operations on types int and integer, and map them to the operations in the target languages.

Contents

Chapter 1

A special case of a conversion.

theory *Code-Int-Integer-Conversion*
imports
  Main
begin

Use this function to convert numeral *integer*s quickly into *int*s. By default, it works only for symbolic evaluation; normally generated code raises an exception at run-time. If theory *Code-Target-Bits-Int* is imported, it works again, because then *int* is implemented in terms of *integer* even for symbolic evaluation.

definition *int-of-integer-symbolic* :: *integer* ⇒ *int*
  where *int-of-integer-symbolic* = *int-of-integer*

lemma *int-of-integer-symbolic-aux-code* [*code nbe*]:
  int-of-integer-symbolic 0 = *0*
  int-of-integer-symbolic (*Code-Numeral.Pos n*) = *Int.Pos n*
  int-of-integer-symbolic (*Code-Numeral.Neg n*) = *Int.Neg n*
  ⟨*proof*⟩

end

Chapter 2

Symbolic implementation of bit operations on int

theory *Code-Symbolic-Bits-Int*
imports
Word-Lib.Least-significant-bit
Word-Lib.Generic-set-bit
Word-Lib.Bit-Comprehension
begin

2.1 Implementations of bit operations on *int* operating on symbolic representation

context
includes *bit-operations-syntax*
begin

lemma *test-bit-int-code* [*code*]:
bit (*0*::*int*) *n* = *False*
bit (*Int.Neg num.One*) *n* = *True*
bit (*Int.Pos num.One*) *0* = *True*
bit (*Int.Pos* (*num.Bit0 m*)) *0* = *False*
bit (*Int.Pos* (*num.Bit1 m*)) *0* = *True*
bit (*Int.Neg* (*num.Bit0 m*)) *0* = *False*
bit (*Int.Neg* (*num.Bit1 m*)) *0* = *True*
bit (*Int.Pos num.One*) (*Suc n*) = *False*
bit (*Int.Pos* (*num.Bit0 m*)) (*Suc n*) = *bit* (*Int.Pos m*) *n*
bit (*Int.Pos* (*num.Bit1 m*)) (*Suc n*) = *bit* (*Int.Pos m*) *n*
bit (*Int.Neg* (*num.Bit0 m*)) (*Suc n*) = *bit* (*Int.Neg m*) *n*
bit (*Int.Neg* (*num.Bit1 m*)) (*Suc n*) = *bit* (*Int.Neg* (*Num.inc m*)) *n*
⟨*proof*⟩

lemma *int-not-code* [*code*]:
NOT (*0* :: *int*) = −*1*

NOT (Int.Pos n) = Int.Neg (Num.inc n)
NOT (Int.Neg n) = Num.sub n num.One
⟨*proof*⟩

lemma *int-and-code* [*code*]: **fixes** *i j* :: *int* **shows**
0 AND j = 0
i AND 0 = 0
Int.Pos n AND Int.Pos m = (case and-num n m of None ⇒ 0 | Some n′ ⇒ Int.Pos n′)
Int.Neg n AND Int.Neg m = NOT (Num.sub n num.One OR Num.sub m num.One)
Int.Pos n AND Int.Neg num.One = Int.Pos n
Int.Pos n AND Int.Neg (num.Bit0 m) = Num.sub (or-not-num-neg (Num.BitM m) n) num.One
Int.Pos n AND Int.Neg (num.Bit1 m) = Num.sub (or-not-num-neg (num.Bit0 m) n) num.One
Int.Neg num.One AND Int.Pos m = Int.Pos m
Int.Neg (num.Bit0 n) AND Int.Pos m = Num.sub (or-not-num-neg (Num.BitM n) m) num.One
Int.Neg (num.Bit1 n) AND Int.Pos m = Num.sub (or-not-num-neg (num.Bit0 n) m) num.One
⟨*proof*⟩

lemma *int-or-code* [*code*]: **fixes** *i j* :: *int* **shows**
0 OR j = j
i OR 0 = i
Int.Pos n OR Int.Pos m = Int.Pos (or-num n m)
Int.Neg n OR Int.Neg m = NOT (Num.sub n num.One AND Num.sub m num.One)
Int.Pos n OR Int.Neg num.One = Int.Neg num.One
Int.Pos n OR Int.Neg (num.Bit0 m) = (case and-not-num (Num.BitM m) n of None ⇒ −1 | Some n′ ⇒ Int.Neg (Num.inc n′))
Int.Pos n OR Int.Neg (num.Bit1 m) = (case and-not-num (num.Bit0 m) n of None ⇒ −1 | Some n′ ⇒ Int.Neg (Num.inc n′))
Int.Neg num.One OR Int.Pos m = Int.Neg num.One
Int.Neg (num.Bit0 n) OR Int.Pos m = (case and-not-num (Num.BitM n) m of None ⇒ −1 | Some n′ ⇒ Int.Neg (Num.inc n′))
Int.Neg (num.Bit1 n) OR Int.Pos m = (case and-not-num (num.Bit0 n) m of None ⇒ −1 | Some n′ ⇒ Int.Neg (Num.inc n′))
⟨*proof*⟩

lemma *int-xor-code* [*code*]: **fixes** *i j* :: *int* **shows**
0 XOR j = j
i XOR 0 = i
Int.Pos n XOR Int.Pos m = (case xor-num n m of None ⇒ 0 | Some n′ ⇒ Int.Pos n′)
Int.Neg n XOR Int.Neg m = Num.sub n num.One XOR Num.sub m num.One
Int.Neg n XOR Int.Pos m = NOT (Num.sub n num.One XOR Int.Pos m)
Int.Pos n XOR Int.Neg m = NOT (Int.Pos n XOR Num.sub m num.One)
⟨*proof*⟩

lemma *bin-rest-code*: *i div 2 = drop-bit 1 i* **for** *i* :: *int*
⟨*proof*⟩

lemma *set-bits-code* [*code*]:
set-bits = Code.abort (*STR ''set-bits is unsupported on type int''*) (λ-. *set-bits* :: - ⇒ *int*)
⟨*proof*⟩

lemma fixes *i* :: *int*
shows *int-set-bit-True-conv-OR* [*code*]: *Generic-set-bit.set-bit i n True = i OR push-bit n 1*
and *int-set-bit-False-conv-NAND* [*code*]: *Generic-set-bit.set-bit i n False = i AND NOT* (*push-bit n 1*)
and *int-set-bit-conv-ops*: *Generic-set-bit.set-bit i n b* = (*if b then i OR* (*push-bit n 1*) *else i AND NOT* (*push-bit n 1*))
⟨*proof*⟩

declare [[*code drop*: ‹*drop-bit* :: *nat* ⇒ *int* ⇒ *int*›]]

lemma *drop-bit-int-code* [*code*]: **fixes** *i* :: *int* **shows**
drop-bit 0 i = i
drop-bit (*Suc n*) *0* = (*0* :: *int*)
drop-bit (*Suc n*) (*Int.Pos num.One*) = *0*
drop-bit (*Suc n*) (*Int.Pos* (*num.Bit0 m*)) = *drop-bit n* (*Int.Pos m*)
drop-bit (*Suc n*) (*Int.Pos* (*num.Bit1 m*)) = *drop-bit n* (*Int.Pos m*)
drop-bit (*Suc n*) (*Int.Neg num.One*) = − *1*
drop-bit (*Suc n*) (*Int.Neg* (*num.Bit0 m*)) = *drop-bit n* (*Int.Neg m*)
drop-bit (*Suc n*) (*Int.Neg* (*num.Bit1 m*)) = *drop-bit n* (*Int.Neg* (*Num.inc m*))
⟨*proof*⟩

declare [[*code drop*: ‹*push-bit* :: *nat* ⇒ *int* ⇒ *int*›]]

lemma *push-bit-int-code* [*code*]:
push-bit 0 i = i
push-bit (*Suc n*) *i = push-bit n* (*Int.dup i*)
⟨*proof*⟩

lemma *int-lsb-code* [*code*]:
lsb (*0* :: *int*) = *False*
lsb (*Int.Pos num.One*) = *True*
lsb (*Int.Pos* (*num.Bit0 w*)) = *False*
lsb (*Int.Pos* (*num.Bit1 w*)) = *True*
lsb (*Int.Neg num.One*) = *True*
lsb (*Int.Neg* (*num.Bit0 w*)) = *False*
lsb (*Int.Neg* (*num.Bit1 w*)) = *True*
⟨*proof*⟩

end

end

Chapter 3

Bit operations for target language integers

theory *Bits-Integer* **imports**
  Word-Lib.Bit-Comprehension
  Code-Int-Integer-Conversion
  Code-Symbolic-Bits-Int
begin

lemmas [*transfer-rule*] =
  identity-quotient
  fun-quotient
  Quotient-integer[*folded integer.pcr-cr-eq*]

lemma *undefined-transfer*:
  assumes *Quotient R Abs Rep T*
  shows *T* (*Rep undefined*) *undefined*
⟨*proof*⟩

bundle *undefined-transfer* = *undefined-transfer*[*transfer-rule*]

3.1 More lemmas about *integer*s

context
includes *integer.lifting*
begin

lemma *bitval-integer-transfer* [*transfer-rule*]:
  (*rel-fun* (=) *pcr-integer*) *of-bool of-bool*
⟨*proof*⟩

lemma *integer-of-nat-less-0-conv* [*simp*]: ¬ *integer-of-nat n* < *0*
⟨*proof*⟩

lemma *int-of-integer-pow*: *int-of-integer* $(x \mathbin{\hat{}} n) = \mathit{int\text{-}of\text{-}integer}\ x \mathbin{\hat{}} n$
⟨*proof*⟩

lemma *pow-integer-transfer* [*transfer-rule*]:
$(\mathit{rel\text{-}fun}\ \mathit{pcr\text{-}integer}\ (\mathit{rel\text{-}fun}\ (=)\ \mathit{pcr\text{-}integer}))\ (\mathbin{\hat{}})\ (\mathbin{\hat{}})$
⟨*proof*⟩

lemma *sub1-lt-0-iff* [*simp*]: $\mathit{Code\text{-}Numeral.sub}\ n\ \mathit{num.One} < 0 \longleftrightarrow \mathit{False}$
⟨*proof*⟩

lemma *nat-of-integer-numeral* [*simp*]: $\mathit{nat\text{-}of\text{-}integer}\ (\mathit{numeral}\ n) = \mathit{numeral}\ n$
⟨*proof*⟩

lemma *nat-of-integer-sub1-conv-pred-numeral* [*simp*]:
$\mathit{nat\text{-}of\text{-}integer}\ (\mathit{Code\text{-}Numeral.sub}\ n\ \mathit{num.One}) = \mathit{pred\text{-}numeral}\ n$
⟨*proof*⟩

lemma *nat-of-integer-1* [*simp*]: $\mathit{nat\text{-}of\text{-}integer}\ 1 = 1$
⟨*proof*⟩

lemma *dup-1* [*simp*]: $\mathit{Code\text{-}Numeral.dup}\ 1 = 2$
⟨*proof*⟩

3.2 Bit operations on *integer*

Bit operations on *integer* are the same as on *int*

lift-definition *bin-rest-integer* :: *integer* ⇒ *integer* **is** ‹λ*k* . *k div* 2› ⟨*proof*⟩
lift-definition *bin-last-integer* :: *integer* ⇒ *bool* **is** *odd* ⟨*proof*⟩
lift-definition *Bit-integer* :: *integer* ⇒ *bool* ⇒ *integer* **is** ‹λ*k b*. *of-bool b* + 2 ∗ *k*›
⟨*proof*⟩

end

instantiation *integer* :: *lsb* **begin**
context includes *integer.lifting* **begin**

lift-definition *lsb-integer* :: *integer* ⇒ *bool* **is** *lsb* ⟨*proof*⟩

instance
⟨*proof*⟩

end
end

instantiation *integer* :: *msb* **begin**
context includes *integer.lifting* **begin**

lift-definition *msb-integer* :: *integer* ⇒ *bool* **is** *msb* ⟨*proof*⟩

instance ⟨*proof*⟩

end
end

instantiation *integer* :: *set-bit* **begin**
context includes *integer.lifting* **begin**

lift-definition *set-bit-integer* :: *integer* ⇒ *nat* ⇒ *bool* ⇒ *integer* **is** *set-bit* ⟨*proof*⟩

instance
⟨*proof*⟩

end
end

abbreviation (*input*) *wf-set-bits-integer*
where *wf-set-bits-integer* ≡ *wf-set-bits-int*

3.3 Target language implementations

Unfortunately, this is not straightforward, because these API functions have different signatures and preconditions on the parameters:

Standard ML Shifts in IntInf are given as word, but not IntInf.

Haskell In the Data.Bits.Bits type class, shifts and bit indices are given as Int rather than Integer.

Additional constants take only parameters of type *integer* rather than *nat* and check the preconditions as far as possible (e.g., being non-negative) in a portable way. Manual implementations inside code_printing perform the remaining range checks and convert these *integer*s into the right type.

For normalisation by evaluation, we derive custom code equations, because NBE does not know these code_printing serialisations and would otherwise loop.

code-identifier code-module *Bits-Integer* ⇀
(*SML*) *Bits-Int* **and** (*OCaml*) *Bits-Int* **and** (*Haskell*) *Bits-Int* **and** (*Scala*) *Bits-Int*

code-printing code-module *Bits-Integer* ⇀ (*SML*)
‹*structure Bits-Integer* : *sig*
val set-bit : *IntInf.int* −> *IntInf.int* −> *bool* −> *IntInf.int*
val shiftl : *IntInf.int* −> *IntInf.int* −> *IntInf.int*
val shiftr : *IntInf.int* −> *IntInf.int* −> *IntInf.int*
val test-bit : *IntInf.int* −> *IntInf.int* −> *bool*

```
end = struct

val maxWord = IntInf.pow (2, Word.wordSize);

fun set-bit x n b =
  if n < maxWord then
    if b then IntInf.orb (x, IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge n)))
   else IntInf.andb (x, IntInf.notb (IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge
n))))
  else raise (Fail (Bit index too large: ^ IntInf.toString n));

fun shiftl x n =
  if n < maxWord then IntInf.<< (x, Word.fromLargeInt (IntInf.toLarge n))
  else raise (Fail (Shift operand too large: ^ IntInf.toString n));

fun shiftr x n =
  if n < maxWord then IntInf.~>> (x, Word.fromLargeInt (IntInf.toLarge n))
  else raise (Fail (Shift operand too large: ^ IntInf.toString n));

fun test-bit x n =
 if n < maxWord then IntInf.andb (x, IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge
n))) <> 0
  else raise (Fail (Bit index too large: ^ IntInf.toString n));

end; (*struct Bits-Integer*)›
code-reserved SML Bits-Integer

code-printing code-module Bits-Integer → (OCaml)
‹module Bits-Integer : sig
  val shiftl : Z.t -> Z.t -> Z.t
  val shiftr : Z.t -> Z.t -> Z.t
  val test-bit : Z.t -> Z.t -> bool
end = struct

(* We do not need an explicit range checks here,
   because Big-int.int-of-big-int raises Failure
   if the argument does not fit into an int. *)
let shiftl x n = Z.shift-left x (Z.to-int n);;

let shiftr x n = Z.shift-right x (Z.to-int n);;

let test-bit x n =  Z.testbit x (Z.to-int n);;

end;; (*struct Bits-Integer*)›
code-reserved OCaml Bits-Integer

code-printing code-module Data-Bits → (Haskell)
‹
module Data-Bits where {
```

```
import qualified Data.Bits;

{-
  The ...Bounded functions assume that the Integer argument for the shift
  or bit index fits into an Int, is non-negative and (for types of fixed bit width)
  less than bitSize
-}

infixl 7 .&.;
infixl 6 `xor`;
infixl 5 .|.;

(.&.) :: Data.Bits.Bits a => a -> a -> a;
(.&.) = (Data.Bits..&.);

xor :: Data.Bits.Bits a => a -> a -> a;
xor = Data.Bits.xor;

(.|.) :: Data.Bits.Bits a => a -> a -> a;
(.|.) = (Data.Bits..|.);

complement :: Data.Bits.Bits a => a -> a;
complement = Data.Bits.complement;

testBitUnbounded :: Data.Bits.Bits a => a -> Integer -> Bool;
testBitUnbounded x b
  | b <= toInteger (Prelude.maxBound :: Int) = Data.Bits.testBit x (fromInteger
b)
  | otherwise = error (Bit index too large: ++ show b)
;

testBitBounded :: Data.Bits.Bits a => a -> Integer -> Bool;
testBitBounded x b = Data.Bits.testBit x (fromInteger b);

setBitUnbounded :: Data.Bits.Bits a => a -> Integer -> Bool -> a;
setBitUnbounded x n b
  | n <= toInteger (Prelude.maxBound :: Int) =
   if b then Data.Bits.setBit x (fromInteger n) else Data.Bits.clearBit x (fromInteger
n)
  | otherwise = error (Bit index too large: ++ show n)
;

setBitBounded :: Data.Bits.Bits a => a -> Integer -> Bool -> a;
setBitBounded x n True = Data.Bits.setBit x (fromInteger n);
setBitBounded x n False = Data.Bits.clearBit x (fromInteger n);

shiftlUnbounded :: Data.Bits.Bits a => a -> Integer -> a;
shiftlUnbounded x n
```

```
  | n <= toInteger (Prelude.maxBound :: Int) = Data.Bits.shiftL x (fromInteger
n)
  | otherwise = error (Shift operand too large: ++ show n)
;

shiftlBounded :: Data.Bits.Bits a => a -> Integer -> a;
shiftlBounded x n = Data.Bits.shiftL x (fromInteger n);

shiftrUnbounded :: Data.Bits.Bits a => a -> Integer -> a;
shiftrUnbounded x n
  | n <= toInteger (Prelude.maxBound :: Int) = Data.Bits.shiftR x (fromInteger
n)
  | otherwise = error (Shift operand too large: ++ show n)
;

shiftrBounded :: (Ord a, Data.Bits.Bits a) => a -> Integer -> a;
shiftrBounded x n = Data.Bits.shiftR x (fromInteger n);

}›
```

and — *HOL.Quickcheck-Narrowing* maps *integer* to Haskell's Prelude.Int type instead of Integer. For compatibility with the Haskell target, we nevertheless provide bounded and unbounded functions.
(*Haskell-Quickcheck*)

```
‹
module Data-Bits where {

import qualified Data.Bits;

{-
  The functions assume that the Int argument for the shift or bit index is
  non-negative and (for types of fixed bit width) less than bitSize
-}

infixl 7 .&.;
infixl 6 `xor`;
infixl 5 .|.;

(.&.) :: Data.Bits.Bits a => a -> a -> a;
(.&.) = (Data.Bits..&.);

xor :: Data.Bits.Bits a => a -> a -> a;
xor = Data.Bits.xor;

(.|.) :: Data.Bits.Bits a => a -> a -> a;
(.|.) = (Data.Bits..|.);

complement :: Data.Bits.Bits a => a -> a;
complement = Data.Bits.complement;
```

```
testBitUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> Bool;
testBitUnbounded = Data.Bits.testBit;

testBitBounded :: Data.Bits.Bits a => a -> Prelude.Int -> Bool;
testBitBounded = Data.Bits.testBit;

setBitUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> Bool -> a;
setBitUnbounded x n True = Data.Bits.setBit x n;
setBitUnbounded x n False = Data.Bits.clearBit x n;

setBitBounded :: Data.Bits.Bits a => a -> Prelude.Int -> Bool -> a;
setBitBounded x n True = Data.Bits.setBit x n;
setBitBounded x n False = Data.Bits.clearBit x n;

shiftlUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> a;
shiftlUnbounded = Data.Bits.shiftL;

shiftlBounded :: Data.Bits.Bits a => a -> Prelude.Int -> a;
shiftlBounded = Data.Bits.shiftL;

shiftrUnbounded :: Data.Bits.Bits a => a -> Prelude.Int -> a;
shiftrUnbounded = Data.Bits.shiftR;

shiftrBounded :: (Ord a, Data.Bits.Bits a) => a -> Prelude.Int -> a;
shiftrBounded = Data.Bits.shiftR;

}›
```

code-reserved *Haskell Data-Bits*

code-printing code-module *Bits-Integer* ⇀ (*Scala*)
‹*object Bits-Integer* {

```
def setBit(x: BigInt, n: BigInt, b: Boolean) : BigInt =
  if (n.isValidInt)
    if (b)
      x.setBit(n.toInt)
    else
      x.clearBit(n.toInt)
  else
    sys.error(Bit index too large: + n.toString)

def shiftl(x: BigInt, n: BigInt) : BigInt =
  if (n.isValidInt)
    x << n.toInt
  else
    sys.error(Shift index too large: + n.toString)

def shiftr(x: BigInt, n: BigInt) : BigInt =
```

```
  if (n.isValidInt)
    x << n.toInt
  else
    sys.error(Shift index too large: + n.toString)

def testBit(x: BigInt, n: BigInt) : Boolean =
  if (n.isValidInt)
    x.testBit(n.toInt)
  else
    sys.error(Bit index too large: + n.toString)

} /* object Bits-Integer */›

code-printing
  constant Bit-Operations.and :: integer ⇒ integer ⇒ integer ⇀
  (SML) IntInf.andb ((-),/ (-)) and
  (OCaml) Z.logand and
  (Haskell) ((Data'-Bits..&.) :: Integer -> Integer -> Integer) and
  (Haskell-Quickcheck) ((Data'-Bits..&.) :: Prelude.Int -> Prelude.Int -> Pre-
lude.Int) and
  (Scala) infixl 3 &
| constant Bit-Operations.or :: integer ⇒ integer ⇒ integer ⇀
  (SML) IntInf.orb ((-),/ (-)) and
  (OCaml) Z.logor and
  (Haskell) ((Data'-Bits..|.) :: Integer -> Integer -> Integer) and
  (Haskell-Quickcheck) ((Data'-Bits..|.) :: Prelude.Int -> Prelude.Int -> Pre-
lude.Int) and
  (Scala) infixl 1 |
| constant Bit-Operations.xor :: integer ⇒ integer ⇒ integer ⇀
  (SML) IntInf.xorb ((-),/ (-)) and
  (OCaml) Z.logxor and
  (Haskell) (Data'-Bits.xor :: Integer -> Integer -> Integer) and
  (Haskell-Quickcheck) (Data'-Bits.xor :: Prelude.Int -> Prelude.Int -> Pre-
lude.Int) and
  (Scala) infixl 2 ^
| constant Bit-Operations.not :: integer ⇒ integer ⇀
  (SML) IntInf.notb and
  (OCaml) Z.lognot and
  (Haskell) (Data'-Bits.complement :: Integer -> Integer) and
 (Haskell-Quickcheck) (Data'-Bits.complement :: Prelude.Int -> Prelude.Int) and
  (Scala) -.unary'-~

code-printing constant bin-rest-integer ⇀
  (SML) IntInf.div ((-), 2) and
  (OCaml) Z.shift'-right/ -/ 1 and
  (Haskell) (Data'-Bits.shiftrUnbounded - 1 :: Integer) and
  (Haskell-Quickcheck) (Data'-Bits.shiftrUnbounded - 1 :: Prelude.Int) and
  (Scala) - >> 1
```

context
includes *integer.lifting bit-operations-syntax*
begin

lemma *bitNOT-integer-code* [*code*]:
NOT i = − i − 1
⟨*proof*⟩

lemma *bin-rest-integer-code* [*code nbe*]:
bin-rest-integer i = i div 2
⟨*proof*⟩

lemma *bin-last-integer-code* [*code*]:
bin-last-integer i ⟷ i AND 1 ≠ 0
⟨*proof*⟩

lemma *bin-last-integer-nbe* [*code nbe*]:
bin-last-integer i ⟷ i mod 2 ≠ 0
⟨*proof*⟩

lemma *bitval-bin-last-integer* [*code-unfold*]:
of-bool (*bin-last-integer i*) *= i AND 1*
⟨*proof*⟩

end

definition *integer-test-bit* :: *integer ⇒ integer ⇒ bool*
where *integer-test-bit x n* = (*if n < 0 then undefined x n else bit x* (*nat-of-integer n*))

declare [[*code drop*: ‹*bit* :: *integer ⇒ nat ⇒ bool*›]]

lemma *bit-integer-code* [*code*]:
bit x n ⟷ integer-test-bit x (*integer-of-nat n*)
⟨*proof*⟩

lemma *integer-test-bit-code* [*code*]:
integer-test-bit x (*Code-Numeral.Neg n*) *= undefined x* (*Code-Numeral.Neg n*)
integer-test-bit 0 0 = False
integer-test-bit 0 (*Code-Numeral.Pos n*) *= False*
integer-test-bit (*Code-Numeral.Pos num.One*) *0 = True*
integer-test-bit (*Code-Numeral.Pos* (*num.Bit0 n*)) *0 = False*
integer-test-bit (*Code-Numeral.Pos* (*num.Bit1 n*)) *0 = True*
integer-test-bit (*Code-Numeral.Pos num.One*) (*Code-Numeral.Pos n′*) *= False*
integer-test-bit (*Code-Numeral.Pos* (*num.Bit0 n*)) (*Code-Numeral.Pos n′*) *=*
integer-test-bit (*Code-Numeral.Pos n*) (*Code-Numeral.sub n′ num.One*)
integer-test-bit (*Code-Numeral.Pos* (*num.Bit1 n*)) (*Code-Numeral.Pos n′*) *=*
integer-test-bit (*Code-Numeral.Pos n*) (*Code-Numeral.sub n′ num.One*)

integer-test-bit (*Code-Numeral.Neg num.One*) *0* = *True*
integer-test-bit (*Code-Numeral.Neg* (*num.Bit0 n*)) *0* = *False*
integer-test-bit (*Code-Numeral.Neg* (*num.Bit1 n*)) *0* = *True*
integer-test-bit (*Code-Numeral.Neg num.One*) (*Code-Numeral.Pos n′*) = *True*
integer-test-bit (*Code-Numeral.Neg* (*num.Bit0 n*)) (*Code-Numeral.Pos n′*) =
integer-test-bit (*Code-Numeral.Neg n*) (*Code-Numeral.sub n′ num.One*)
integer-test-bit (*Code-Numeral.Neg* (*num.Bit1 n*)) (*Code-Numeral.Pos n′*) =
integer-test-bit (*Code-Numeral.Neg* (*n* + *num.One*)) (*Code-Numeral.sub n′ num.One*)
⟨*proof*⟩

code-printing constant *integer-test-bit* ⇀
(*SML*) *Bits′-Integer.test′-bit* **and**
(*OCaml*) *Bits′-Integer.test′-bit* **and**
(*Haskell*) (*Data′-Bits.testBitUnbounded* :: *Integer* −> *Integer* −> *Bool*) **and**
(*Haskell-Quickcheck*) (*Data′-Bits.testBitUnbounded* :: *Prelude.Int* −> *Prelude.Int* −> *Bool*) **and**
(*Scala*) *Bits′-Integer.testBit*

context
includes *integer.lifting bit-operations-syntax*
begin

lemma *lsb-integer-code* [*code*]:
fixes *x* :: *integer* **shows**
lsb x = *bit x 0*
⟨*proof*⟩

definition *integer-set-bit* :: *integer* ⇒ *integer* ⇒ *bool* ⇒ *integer*
where [*code del*]: *integer-set-bit x n b* = (*if n* < *0 then undefined x n b else set-bit x* (*nat-of-integer n*) *b*)

lemma *set-bit-integer-code* [*code*]:
set-bit x i b = *integer-set-bit x* (*integer-of-nat i*) *b*
⟨*proof*⟩

lemma *set-bit-integer-conv-masks*:
fixes *x* :: *integer* **shows**
set-bit x i b = (*if b then x OR* (*push-bit i 1*) *else x AND NOT* (*push-bit i 1*))
⟨*proof*⟩

end

code-printing constant *integer-set-bit* ⇀
(*SML*) *Bits′-Integer.set′-bit* **and**
(*Haskell*) (*Data′-Bits.setBitUnbounded* :: *Integer* −> *Integer* −> *Bool* −> *Integer*) **and**
(*Haskell-Quickcheck*) (*Data′-Bits.setBitUnbounded* :: *Prelude.Int* −> *Prelude.Int* −> *Bool* −> *Prelude.Int*) **and**
(*Scala*) *Bits′-Integer.setBit*

OCaml.Big_int does not have a method for changing an individual bit, so we emulate that with masks. We prefer an Isabelle implementation, because this then takes care of the signs for AND and OR.

context
includes *bit-operations-syntax*
begin

lemma *integer-set-bit-code* [*code*]:
integer-set-bit x n b =
(*if n* < *0 then undefined x n b else*
if b then x OR (*push-bit* (*nat-of-integer n*) *1*)
else x AND NOT (*push-bit* (*nat-of-integer n*) *1*))
⟨*proof*⟩

end

definition *integer-shiftl* :: *integer* ⇒ *integer* ⇒ *integer*
where [*code del*]: *integer-shiftl x n* = (*if n* < *0 then undefined x n else push-bit* (*nat-of-integer n*) *x*)

declare [[*code drop*: ‹*push-bit* :: *nat* ⇒ *integer* ⇒ *integer*›]]

lemma *shiftl-integer-code* [*code*]:
fixes *x* :: *integer* **shows**
push-bit n x = *integer-shiftl x* (*integer-of-nat n*)
⟨*proof*⟩

context
includes *integer.lifting*
begin

lemma *shiftl-integer-conv-mult-pow2*:
fixes *x* :: *integer* **shows**
push-bit n x = *x* ∗ *2* ^ *n*
⟨*proof*⟩

lemma *integer-shiftl-code* [*code*]:
integer-shiftl x (*Code-Numeral.Neg n*) = *undefined x* (*Code-Numeral.Neg n*)
integer-shiftl x 0 = *x*
integer-shiftl x (*Code-Numeral.Pos n*) = *integer-shiftl* (*Code-Numeral.dup x*) (*Code-Numeral.sub n num.One*)
integer-shiftl 0 (*Code-Numeral.Pos n*) = *0*
⟨*proof*⟩

end

code-printing constant *integer-shiftl* ⇀
(*SML*) *Bits'-Integer.shiftl* **and**
(*OCaml*) *Bits'-Integer.shiftl* **and**

(*Haskell*) (*Data′-Bits.shiftlUnbounded* :: *Integer* −> *Integer* −> *Integer*) **and**
(*Haskell-Quickcheck*) (*Data′-Bits.shiftlUnbounded* :: *Prelude.Int* −> *Prelude.Int* −> *Prelude.Int*) **and**
(*Scala*) *Bits′-Integer.shiftl*

definition *integer-shiftr* :: *integer* ⇒ *integer* ⇒ *integer*
where [*code del*]: *integer-shiftr x n* = (*if n* < *0 then undefined x n else drop-bit* (*nat-of-integer n*) *x*)

declare [[*code drop*: ‹*drop-bit* :: *nat* ⇒ *integer* ⇒ *integer*›]]

lemma *shiftr-integer-conv-div-pow2*:
includes *integer.lifting* **fixes** *x* :: *integer* **shows**
drop-bit n x = *x div* 2 ^ *n*
⟨*proof*⟩

lemma *shiftr-integer-code* [*code*]:
fixes *x* :: *integer* **shows**
drop-bit n x = *integer-shiftr x* (*integer-of-nat n*)
⟨*proof*⟩

code-printing constant *integer-shiftr* ⇀
(*SML*) *Bits′-Integer.shiftr* **and**
(*OCaml*) *Bits′-Integer.shiftr* **and**
(*Haskell*) (*Data′-Bits.shiftrUnbounded* :: *Integer* −> *Integer* −> *Integer*) **and**
(*Haskell-Quickcheck*) (*Data′-Bits.shiftrUnbounded* :: *Prelude.Int* −> *Prelude.Int* −> *Prelude.Int*) **and**
(*Scala*) *Bits′-Integer.shiftr*

lemma *integer-shiftr-code* [*code*]:
includes *integer.lifting*
shows
integer-shiftr x (*Code-Numeral.Neg n*) = *undefined x* (*Code-Numeral.Neg n*)
integer-shiftr x 0 = *x*
integer-shiftr 0 (*Code-Numeral.Pos n*) = *0*
integer-shiftr (*Code-Numeral.Pos num.One*) (*Code-Numeral.Pos n*) = *0*
integer-shiftr (*Code-Numeral.Pos* (*num.Bit0 n′*)) (*Code-Numeral.Pos n*) =
integer-shiftr (*Code-Numeral.Pos n′*) (*Code-Numeral.sub n num.One*)
integer-shiftr (*Code-Numeral.Pos* (*num.Bit1 n′*)) (*Code-Numeral.Pos n*) =
integer-shiftr (*Code-Numeral.Pos n′*) (*Code-Numeral.sub n num.One*)
integer-shiftr (*Code-Numeral.Neg num.One*) (*Code-Numeral.Pos n*) = −*1*
integer-shiftr (*Code-Numeral.Neg* (*num.Bit0 n′*)) (*Code-Numeral.Pos n*) =
integer-shiftr (*Code-Numeral.Neg n′*) (*Code-Numeral.sub n num.One*)
integer-shiftr (*Code-Numeral.Neg* (*num.Bit1 n′*)) (*Code-Numeral.Pos n*) =
integer-shiftr (*Code-Numeral.Neg* (*Num.inc n′*)) (*Code-Numeral.sub n num.One*)
⟨*proof*⟩

context
includes *integer.lifting*

begin

lemma *Bit-integer-code* [*code*]:
Bit-integer i False = push-bit 1 i
Bit-integer i True = push-bit 1 i + 1
⟨*proof*⟩

lemma *msb-integer-code* [*code*]:
msb (*x* :: *integer*) ⟷ *x* < *0*
⟨*proof*⟩

end

context
includes *integer.lifting natural.lifting bit-operations-syntax*
begin

lemma *bitAND-integer-unfold* [*code*]:
x AND y =
(*if x* = *0 then 0*
else if x = − *1 then y*
else Bit-integer (*bin-rest-integer x AND bin-rest-integer y*) (*bin-last-integer x* ∧ *bin-last-integer y*))
⟨*proof*⟩

lemma *bitOR-integer-unfold* [*code*]:
x OR y =
(*if x* = *0 then y*
else if x = − *1 then* − *1*
else Bit-integer (*bin-rest-integer x OR bin-rest-integer y*) (*bin-last-integer x* ∨ *bin-last-integer y*))
⟨*proof*⟩

lemma *bitXOR-integer-unfold* [*code*]:
x XOR y =
(*if x* = *0 then y*
else if x = − *1 then NOT y*
else Bit-integer (*bin-rest-integer x XOR bin-rest-integer y*)
(¬ *bin-last-integer x* ⟷ *bin-last-integer y*))
⟨*proof*⟩

end

3.4 Test code generator setup

context
includes *bit-operations-syntax*
begin

definition *bit-integer-test* :: *bool* **where**
bit-integer-test =
(([*−1 AND 3*, *1 AND −3*, *3 AND 5*, *−3 AND* (*− 5*)
, *−3 OR 1*, *1 OR −3*, *3 OR 5*, *−3 OR* (*− 5*)
, *NOT 1*, *NOT* (*− 3*)
, *−1 XOR 3*, *1 XOR* (*− 3*), *3 XOR 5*, *−5 XOR* (*− 3*)
, *set-bit 5 4 True*, *set-bit* (*− 5*) *2 True*, *set-bit 5 0 False*, *set-bit* (*− 5*) *1 False*
, *push-bit 2 1*, *push-bit 3* (*− 1*)
, *drop-bit 3 100*, *drop-bit 3* (*− 100*)] :: *integer list*)
= [*3*, *1*, *1*, *−7*
, *−3*, *−3*, *7*, *−1*
, *−2*, *2*
, *−4*, *−4*, *6*, *6*
, *21*, *−1*, *4*, *−7*
, *4*, *−8*
, *12*, *−13*] ∧
[*bit* (*5* :: *integer*) *4*, *bit* (*5* :: *integer*) *2*, *bit* (*−5* :: *integer*) *4*, *bit* (*−5* :: *integer*) *2*
, *lsb* (*5* :: *integer*), *lsb* (*4* :: *integer*), *lsb* (*−1* :: *integer*), *lsb* (*−2* :: *integer*),
msb (*5* :: *integer*), *msb* (*0* :: *integer*), *msb* (*−1* :: *integer*), *msb* (*−2* :: *integer*)]
= [*False*, *True*, *True*, *False*,
True, *False*, *True*, *False*,
False, *False*, *True*, *True*])

export-code *bit-integer-test* **checking** *SML Haskell? Haskell-Quickcheck? OCaml? Scala*

notepad begin
⟨*proof*⟩
end
⟨*ML*⟩

lemma *x AND y* = *x OR* (*y* :: *integer*)
quickcheck[*random*, *expect*=*counterexample*]
quickcheck[*exhaustive*, *expect*=*counterexample*]
⟨*proof*⟩

lemma (*x* :: *integer*) *AND x* = *x OR x*
quickcheck[*narrowing*, *expect*=*no-counterexample*]
⟨*proof*⟩

lemma (*f* :: *integer* ⇒ *unit*) = *g*
quickcheck[*narrowing*, *size*=*3*, *expect*=*no-counterexample*]
⟨*proof*⟩

end

hide-const *bit-integer-test*
hide-fact *bit-integer-test-def*

end

Chapter 4

Common base for target language implementations of word types

theory *Code-Target-Word-Base* **imports**
HOL−Library.Word
Word-Lib.Signed-Division-Word
Bits-Integer
begin

More lemmas

lemma *div-half-nat*:
fixes *x y* :: *nat*
assumes *y* ≠ *0*
shows (*x div y*, *x mod y*) = (*let q* = *2* ∗ (*x div 2 div y*); *r* = *x* − *q* ∗ *y in if y* ≤ *r then* (*q* + *1*, *r* − *y*) *else* (*q*, *r*))
⟨*proof*⟩

lemma *div-half-word*:
fixes *x y* :: *'a* :: *len word*
assumes *y* ≠ *0*
shows (*x div y*, *x mod y*) = (*let q* = *push-bit 1* (*drop-bit 1 x div y*); *r* = *x* − *q* ∗ *y in if y* ≤ *r then* (*q* + *1*, *r* − *y*) *else* (*q*, *r*))
⟨*proof*⟩

lemma *word-test-bit-set-bits*: *bit* (*BITS n. f n* :: *'a* :: *len word*) *n* ⟷ *n* < *LENGTH*(*'a*) ∧ *f n*
⟨*proof*⟩

lemma *word-of-int-conv-set-bits*: *word-of-int i* = (*BITS n. bit i n*)
⟨*proof*⟩

context

includes *bit-operations-syntax*
begin

lemma *word-and-mask-or-conv-and-mask*:
bit n index $\Longrightarrow$ (*n AND mask index*) *OR* (*push-bit index 1*) = *n AND mask* (*index* + *1*)
for *n* :: ‹*'a*::*len word*›
⟨*proof*⟩

lemma *uint-and-mask-or-full*:
fixes *n* :: *'a* :: *len word*
assumes *bit n* (*LENGTH*(*'a*) − *1*)
and *mask1* = *mask* (*LENGTH*(*'a*) − *1*)
and *mask2* = *push-bit* (*LENGTH*(*'a*) − *1*) *1*
shows *uint* (*n AND mask1*) *OR mask2* = *uint n*
⟨*proof*⟩

end

Division on *'a word* is unsigned, but Scala and OCaml only have signed division and modulus.

lemmas *word-sdiv-def* = *sdiv-word-def*
lemmas *word-smod-def* = *smod-word-def*

lemma [*code*]:
x sdiv y =
(*let x'* = *sint x*; *y'* = *sint y*;
negative = (*x'* < *0*) ≠ (*y'* < *0*);
result = *abs x' div abs y'*
in word-of-int (*if negative then* −*result else result*))
for *x y* :: ‹*'a*::*len word*›
⟨*proof*⟩

lemma [*code*]:
x smod y =
(*let x'* = *sint x*; *y'* = *sint y*;
negative = (*x'* < *0*);
result = *abs x' mod abs y'*
in word-of-int (*if negative then* −*result else result*))
for *x y* :: ‹*'a*::*len word*›
⟨*proof*⟩

This algorithm implements unsigned division in terms of signed division. Taken from Hacker's Delight.

lemma *divmod-via-sdivmod*:
fixes *x y* :: *'a* :: *len word*
assumes *y* ≠ *0*
shows
(*x div y*, *x mod y*) =

(*if push-bit* (*LENGTH*('*a*) − *1*) *1* ≤ *y then if x* < *y then* (*0*, *x*) *else* (*1*, *x* − *y*)
else let q = (*push-bit 1* (*drop-bit 1 x sdiv y*));
r = *x* − *q* ∗ *y*
in if r ≥ *y then* (*q* + *1*, *r* − *y*) *else* (*q*, *r*))
⟨*proof*⟩

More implementations tailored towards target-language implementations

context
includes *integer.lifting*
begin

lift-definition *word-of-integer* :: *integer* ⇒ '*a* :: *len word* **is** *word-of-int* ⟨*proof*⟩

lemma *word-of-integer-code* [*code*]: *word-of-integer n* = *word-of-int* (*int-of-integer n*)
⟨*proof*⟩

end

context
includes *bit-operations-syntax*
begin

lemma *word-of-int-code*:
uint (*word-of-int x* :: '*a word*) = *x AND mask* (*LENGTH*('*a* :: *len*))
⟨*proof*⟩

context
fixes *f* :: *nat* ⇒ *bool*
begin

definition *set-bits-aux* :: ‹*nat* ⇒ '*a word* ⇒ '*a*::*len word*›
where ‹*set-bits-aux n w* = *push-bit n w OR take-bit n* (*set-bits f*)›

lemma *bit-set-bit-aux* [*bit-simps*]:
‹*bit* (*set-bits-aux n w*) *m* ⟷ *m* < *LENGTH*('*a*) ∧
(*if m* < *n then f m else bit w* (*m* − *n*))› **for** *w* :: ‹'*a*::*len word*›
⟨*proof*⟩

corollary *set-bits-conv-set-bits-aux*:
‹*set-bits f* = (*set-bits-aux LENGTH*('*a*) *0* :: '*a* :: *len word*)›
⟨*proof*⟩

lemma *set-bits-aux-0* [*simp*]:
‹*set-bits-aux 0 w* = *w*›
⟨*proof*⟩

lemma *set-bits-aux-Suc* [*simp*]:
‹*set-bits-aux* (*Suc n*) *w* = *set-bits-aux n* (*push-bit 1 w OR* (*if f n then 1 else 0*))›

⟨*proof*⟩

lemma *set-bits-aux-simps* [*code*]:
‹*set-bits-aux 0 w = w*›
‹*set-bits-aux (Suc n) w = set-bits-aux n (push-bit 1 w OR (if f n then 1 else 0))*›
⟨*proof*⟩

lemma *set-bits-aux-rec*:
‹*set-bits-aux n w =*
(if n = 0 then w
else let n′ = n − 1 in set-bits-aux n′ (push-bit 1 w OR (if f n′ then 1 else 0)))›
⟨*proof*⟩

end

lemma *word-of-int-via-signed*:
fixes *mask*
assumes *mask-def*: *mask = Bit-Operations.mask (LENGTH('a))*
and *shift-def*: *shift = push-bit LENGTH('a) 1*
and *index-def*: *index = LENGTH('a) − 1*
and *overflow-def*:*overflow = push-bit (LENGTH('a) − 1) 1*
and *least-def*: *least = − overflow*
shows
(word-of-int i :: 'a :: len word) =
(let i′ = i AND mask
in if bit i′ index then
if i′ − shift < least ∨ overflow ≤ i′ − shift then arbitrary1 i′ else word-of-int
(i′ − shift)
else if i′ < least ∨ overflow ≤ i′ then arbitrary2 i′ else word-of-int i′)
⟨*proof*⟩

end

Quickcheck conversion functions

context
includes *state-combinator-syntax*
begin

definition *qc-random-cnv* ::
(natural ⇒ 'a::term-of) ⇒ natural ⇒ Random.seed
⇒ ('a × (unit ⇒ Code-Evaluation.term)) × Random.seed
where *qc-random-cnv a-of-natural i = Random.range (i + 1) ∘→ (λk. Pair (*
let n = a-of-natural k
in (n, λ-. Code-Evaluation.term-of n)))

end

definition *qc-exhaustive-cnv* :: *(natural ⇒ 'a) ⇒ ('a ⇒ (bool × term list) option)*
⇒ natural ⇒ (bool × term list) option

where
qc-exhaustive-cnv a-of-natural f d =
Quickcheck-Exhaustive.exhaustive (%x. f (a-of-natural x)) d

definition *qc-full-exhaustive-cnv* ::
(natural ⇒ ('a::term-of)) ⇒ ('a × (unit ⇒ term) ⇒ (bool × term list) option)
⇒ natural ⇒ (bool × term list) option
where
qc-full-exhaustive-cnv a-of-natural f d = Quickcheck-Exhaustive.full-exhaustive
(%(x, xt). f (a-of-natural x, %-. Code-Evaluation.term-of (a-of-natural x))) d

declare [[*quickcheck-narrowing-ghc-options* = *−XTypeSynonymInstances*]]

definition *qc-narrowing-drawn-from* :: *'a list ⇒ integer ⇒* -
where
qc-narrowing-drawn-from xs =
foldr Quickcheck-Narrowing.sum (map Quickcheck-Narrowing.cons (butlast xs))
(Quickcheck-Narrowing.cons (last xs))

locale *quickcheck-narrowing-samples =*
fixes *a-of-integer* :: *integer ⇒ 'a × 'a* :: {*partial-term-of*, *term-of*}
and *zero* :: *'a*
and *tr* :: *typerep*
begin

function *narrowing-samples* :: *integer ⇒ 'a list*
where
narrowing-samples i =
(if i > 0 then let (a, a') = a-of-integer i in narrowing-samples (i − 1) @ [a, a']
else [zero])
⟨*proof*⟩
termination including *integer.lifting*
⟨*proof*⟩

definition *partial-term-of-sample* :: *integer ⇒ 'a*
where
partial-term-of-sample i =
(if i < 0 then undefined
else if i = 0 then zero
else if i mod 2 = 0 then snd (a-of-integer (i div 2))
else fst (a-of-integer (i div 2 + 1)))

lemma *partial-term-of-code*:
partial-term-of (ty :: 'a itself) (Quickcheck-Narrowing.Narrowing-variable p t) ≡
Code-Evaluation.Free (STR ''-'') tr
partial-term-of (ty :: 'a itself) (Quickcheck-Narrowing.Narrowing-constructor i
[]) ≡
Code-Evaluation.term-of (partial-term-of-sample i)
⟨*proof*⟩

end

lemmas [*code*] =
quickcheck-narrowing-samples.narrowing-samples.simps
quickcheck-narrowing-samples.partial-term-of-sample-def

The separate code target *SML-word* collects setups for the code generator that PolyML does not provide.

⟨*ML*⟩

code-identifier code-module *Code-Target-Word-Base* ⇀
(*SML*) *Word* **and** (*Haskell*) *Word* **and** (*OCaml*) *Word* **and** (*Scala*) *Word*

end

Chapter 5

Systematic approach towards type copies of word type

theory *Word-Type-Copies*
  imports
    HOL−Library.Word
    Word-Lib.Most-significant-bit
    Word-Lib.Least-significant-bit
    Word-Lib.Generic-set-bit
    Code-Target-Word-Base
begin

lemma *int-of-integer-unsigned-eq* [*simp*]:
  ‹*int-of-integer* (*unsigned w*) = *uint w*›
  ⟨*proof*⟩

lemma *int-of-integer-signed-eq* [*simp*]:
  ‹*int-of-integer* (*signed w*) = *sint w*›
  ⟨*proof*⟩

lemma *word-of-int-of-integer-eq* [*simp*]:
  ‹*word-of-int* (*int-of-integer k*) = *word-of-integer k*›
  ⟨*proof*⟩

The lifting machinery is not localized, hence the abstract proofs are carried out using morphisms.

locale *word-type-copy* =
  fixes *of-word* :: ‹'*b*::*len word* ⇒ '*a*›
    and *word-of* :: ‹'*a* ⇒ '*b word*›
  assumes *type-definition*: ‹*type-definition word-of of-word UNIV*›
begin

lemma *word-of-word*:
  ‹*word-of* (*of-word w*) = *w*›
  ⟨*proof*⟩

lemma *of-word-of* [*code abstype*]:
‹*of-word* (*word-of p*) = *p*›
— Use an abstract type for code generation to disable pattern matching on *of-word*.
⟨*proof*⟩

lemma *word-of-eqI*:
‹*p* = *q*› **if** ‹*word-of p* = *word-of q*›
⟨*proof*⟩

lemma *eq-iff-word-of*:
‹*p* = *q* ⟷ *word-of p* = *word-of q*›
⟨*proof*⟩

end

bundle *constraintless*
begin

⟨*ML*⟩

end

locale *word-type-copy-ring* = *word-type-copy*
opening *constraintless* +
constrains *word-of* :: ‹'*a* ⇒ '*b*::*len word*›
assumes *word-of-0* [*code*]: ‹*word-of 0* = *0*›
and *word-of-1* [*code*]: ‹*word-of 1* = *1*›
and *word-of-add* [*code*]: ‹*word-of* (*p* + *q*) = *word-of p* + *word-of q*›
and *word-of-minus* [*code*]: ‹*word-of* (− *p*) = − (*word-of p*)›
and *word-of-diff* [*code*]: ‹*word-of* (*p* − *q*) = *word-of p* − *word-of q*›
and *word-of-mult* [*code*]: ‹*word-of* (*p* ∗ *q*) = *word-of p* ∗ *word-of q*›
and *word-of-div* [*code*]: ‹*word-of* (*p div q*) = *word-of p div word-of q*›
and *word-of-mod* [*code*]: ‹*word-of* (*p mod q*) = *word-of p mod word-of q*›
and *equal-iff-word-of* [*code*]: ‹*HOL.equal p q* ⟷ *HOL.equal* (*word-of p*) (*word-of q*)›
and *less-eq-iff-word-of* [*code*]: ‹*p* ≤ *q* ⟷ *word-of p* ≤ *word-of q*›
and *less-iff-word-of* [*code*]: ‹*p* < *q* ⟷ *word-of p* < *word-of q*›
begin

lemma *of-class-comm-ring-1*:
‹*OFCLASS*('*a*, *comm-ring-1-class*)›
⟨*proof*⟩

lemma *of-class-semiring-modulo*:
‹*OFCLASS*('*a*, *semiring-modulo-class*)›
⟨*proof*⟩

lemma *of-class-equal*:
‹*OFCLASS*('*a*, *equal-class*)›
⟨*proof*⟩

lemma *of-class-linorder*:
‹*OFCLASS*('*a*, *linorder-class*)›
⟨*proof*⟩

end

locale *word-type-copy-bits* = *word-type-copy-ring*
opening *constraintless bit-operations-syntax* +
constrains *word-of* :: ‹'*a*::{*comm-ring-1*, *semiring-modulo*, *equal*, *linorder*} ⇒ '*b*::*len word*›
fixes *signed-drop-bit* :: ‹*nat* ⇒ '*a* ⇒ '*a*›
assumes *bit-eq-word-of* [*code*]: ‹*bit p* = *bit* (*word-of p*)›
and *word-of-not* [*code*]: ‹*word-of* (*NOT p*) = *NOT* (*word-of p*)›
and *word-of-and* [*code*]: ‹*word-of* (*p AND q*) = *word-of p AND word-of q*›
and *word-of-or* [*code*]: ‹*word-of* (*p OR q*) = *word-of p OR word-of q*›
and *word-of-xor* [*code*]: ‹*word-of* (*p XOR q*) = *word-of p XOR word-of q*›
and *word-of-mask* [*code*]: ‹*word-of* (*mask n*) = *mask n*›
and *word-of-push-bit* [*code*]: ‹*word-of* (*push-bit n p*) = *push-bit n* (*word-of p*)›
and *word-of-drop-bit* [*code*]: ‹*word-of* (*drop-bit n p*) = *drop-bit n* (*word-of p*)›
and *word-of-signed-drop-bit* [*code*]: ‹*word-of* (*signed-drop-bit n p*) = *Word.signed-drop-bit n* (*word-of p*)›
and *word-of-take-bit* [*code*]: ‹*word-of* (*take-bit n p*) = *take-bit n* (*word-of p*)›
and *word-of-set-bit* [*code*]: ‹*word-of* (*Bit-Operations.set-bit n p*) = *Bit-Operations.set-bit n* (*word-of p*)›
and *word-of-unset-bit* [*code*]: ‹*word-of* (*unset-bit n p*) = *unset-bit n* (*word-of p*)›
and *word-of-flip-bit* [*code*]: ‹*word-of* (*flip-bit n p*) = *flip-bit n* (*word-of p*)›
begin

lemma *word-of-bool*:
‹*word-of* (*of-bool n*) = *of-bool n*›
⟨*proof*⟩

lemma *word-of-nat*:
‹*word-of* (*of-nat n*) = *of-nat n*›
⟨*proof*⟩

lemma *word-of-numeral* [*simp*]:
‹*word-of* (*numeral n*) = *numeral n*›
⟨*proof*⟩

lemma *word-of-power*:
‹*word-of* (*p* ^ *n*) = *word-of p* ^ *n*›
⟨*proof*⟩

lemma *even-iff-word-of*:
‹*2 dvd p* ⟷ *2 dvd word-of p*› (**is** ‹*?P* ⟷ *?Q*›)
⟨*proof*⟩

lemma *of-class-ring-bit-operations*:
‹*OFCLASS*('*a*, *ring-bit-operations-class*)›
⟨*proof*⟩

lemma [*code*]:
‹*take-bit n a* = *a AND mask n*› **for** *a* :: '*a*
⟨*proof*⟩

lemma [*code*]:
‹*mask* (*Suc n*) = *push-bit n* (*1* :: '*a*) *OR mask n*›
‹*mask 0* = (*0* :: '*a*)›
⟨*proof*⟩

lemma [*code*]:
‹*Bit-Operations.set-bit n w* = *w OR push-bit n 1*› **for** *w* :: '*a*
⟨*proof*⟩

lemma [*code*]:
‹*unset-bit n w* = *w AND NOT* (*push-bit n 1*)› **for** *w* :: '*a*
⟨*proof*⟩

lemma [*code*]:
‹*flip-bit n w* = *w XOR push-bit n 1*› **for** *w* :: '*a*
⟨*proof*⟩

end

locale *word-type-copy-more* = *word-type-copy-bits* +
constrains *word-of* :: ‹'*a*::{*ring-bit-operations*, *equal*, *linorder*} ⇒ '*b*::*len word*›
fixes *of-nat* :: ‹*nat* ⇒ '*a*› **and** *nat-of* :: ‹'*a* ⇒ *nat*›
and *of-int* :: ‹*int* ⇒ '*a*› **and** *int-of* :: ‹'*a* ⇒ *int*›
and *of-integer* :: ‹*integer* ⇒ '*a*› **and** *integer-of* :: ‹'*a* ⇒ *integer*›
assumes *word-of-nat-eq*: ‹*word-of* (*of-nat n*) = *word-of-nat n*›
and *nat-of-eq-word-of*: ‹*nat-of p* = *unat* (*word-of p*)›
and *word-of-int-eq*: ‹*word-of* (*of-int k*) = *word-of-int k*›
and *int-of-eq-word-of*: ‹*int-of p* = *uint* (*word-of p*)›
and *word-of-integer-eq*: ‹*word-of* (*of-integer l*) = *word-of-integer l*›
and *integer-of-eq-word-of*: ‹*integer-of p* = *unsigned* (*word-of p*)›
begin

lemma *of-word-numeral* [*code-post*]:
‹*of-word* (*numeral n*) = *numeral n*›
⟨*proof*⟩

lemma *of-word-0* [*code-post*]:

‹*of-word 0 = 0*›
⟨*proof*⟩

lemma *of-word-1* [*code-post*]:
‹*of-word 1 = 1*›
⟨*proof*⟩

Use pretty numerals from integer for pretty printing

lemma *numeral-eq-integer* [*code-unfold*]:
‹*numeral n = of-integer (numeral n)*›
⟨*proof*⟩

lemma *neg-numeral-eq-integer* [*code-unfold*]:
‹*− numeral n = of-integer (− numeral n)*›
⟨*proof*⟩

end

locale *word-type-copy-misc = word-type-copy-more*
opening *constraintless bit-operations-syntax +*
constrains *word-of* :: ‹*'a*::{*ring-bit-operations*, *equal*, *linorder*} ⇒ *'b*::*len word*›
fixes *size* :: *nat* **and** *set-bits-aux* :: ‹(*nat* ⇒ *bool*) ⇒ *nat* ⇒ *'a* ⇒ *'a*›
assumes *size-eq-length*: ‹*size = LENGTH('b*::*len)›
and *msb-iff-word-of* [*code*]: ‹*msb p ⟷ msb (word-of p)*›
and *lsb-iff-word-of* [*code*]: ‹*lsb p ⟷ lsb (word-of p)*›
and *size-eq-word-of*: ‹*Nat.size (p :: 'a) = Nat.size (word-of p)*›
and *word-of-generic-set-bit* [*code*]: ‹*word-of (Generic-set-bit.set-bit p n b) =*
Generic-set-bit.set-bit (word-of p) n b›
and *word-of-set-bits*: ‹*word-of (set-bits P) = set-bits P*›
and *word-of-set-bits-aux*: ‹*word-of (set-bits-aux P n p) = Code-Target-Word-Base.set-bits-aux*
P n (word-of p)›
begin

lemma *size-eq* [*code*]:
‹*Nat.size p = size*› **for** *p* :: *'a*
⟨*proof*⟩

lemma *set-bits-aux-code* [*code*]:
‹*set-bits-aux f n w =*
(*if n = 0 then w*
else let n' = n − 1 in set-bits-aux f n' (push-bit 1 w OR (if f n' then 1 else 0)))›
⟨*proof*⟩

lemma *set-bits-code* [*code*]: ‹*set-bits P = set-bits-aux P size 0*›
⟨*proof*⟩

lemma *of-class-lsb*:
‹*OFCLASS('a, lsb-class)*›
⟨*proof*⟩

lemma *of-class-set-bit*:
‹*OFCLASS*('*a*, *set-bit-class*)›
⟨*proof*⟩

lemma *of-class-bit-comprehension*:
‹*OFCLASS*('*a*, *bit-comprehension-class*)›
⟨*proof*⟩

end

end

Chapter 6

Unsigned words of 64 bits

theory *Uint64* **imports**
Code-Target-Word-Base Word-Type-Copies
begin

PolyML (in version 5.7) provides a Word64 structure only when run in 64-bit mode. Therefore, we by default provide an implementation of 64-bit words using `IntInf.int` and masking. The code target `SML_word` replaces this implementation and maps the operations directly to the `Word64` structure provided by the Standard ML implementations.

The `Eval` target used by **value** and *eval* dynamically tests at runtime for the version of PolyML and uses PolyML's Word64 structure if it detects a 64-bit version which does not suffer from a division bug found in PolyML 5.6.

6.1 Type definition and primitive operations

typedef *uint64* = ‹*UNIV* :: *64 word set*› ⟨*proof*⟩

global-interpretation *uint64*: *word-type-copy Abs-uint64 Rep-uint64*
⟨*proof*⟩

setup-lifting *type-definition-uint64*

declare *uint64.of-word-of* [*code abstype*]

declare *Quotient-uint64* [*transfer-rule*]

instantiation *uint64* :: ‹{*comm-ring-1*, *semiring-modulo*, *equal*, *linorder*}›
begin

lift-definition *zero-uint64* :: *uint64* **is** *0* ⟨*proof*⟩
lift-definition *one-uint64* :: *uint64* **is** *1* ⟨*proof*⟩
lift-definition *plus-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹(+)› ⟨*proof*⟩

lift-definition *uminus-uint64* :: ‹*uint64* ⇒ *uint64*› **is** *uminus* ⟨*proof*⟩
lift-definition *minus-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹(−)› ⟨*proof*⟩
lift-definition *times-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹(∗)› ⟨*proof*⟩
lift-definition *divide-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹(*div*)› ⟨*proof*⟩
lift-definition *modulo-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹(*mod*)› ⟨*proof*⟩
lift-definition *equal-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *bool*› **is** ‹*HOL.equal*› ⟨*proof*⟩
lift-definition *less-eq-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *bool*› **is** ‹(≤)› ⟨*proof*⟩
lift-definition *less-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *bool*› **is** ‹(<)› ⟨*proof*⟩

global-interpretation *uint64*: *word-type-copy-ring Abs-uint64 Rep-uint64*
⟨*proof*⟩

instance ⟨*proof*⟩

end

instantiation *uint64* :: *ring-bit-operations*
begin

lift-definition *bit-uint64* :: ‹*uint64* ⇒ *nat* ⇒ *bool*› **is** *bit* ⟨*proof*⟩
lift-definition *not-uint64* :: ‹*uint64* ⇒ *uint64*› **is** ‹*Bit-Operations.not*› ⟨*proof*⟩
lift-definition *and-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹*Bit-Operations.and*›
⟨*proof*⟩
lift-definition *or-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹*Bit-Operations.or*›
⟨*proof*⟩
lift-definition *xor-uint64* :: ‹*uint64* ⇒ *uint64* ⇒ *uint64*› **is** ‹*Bit-Operations.xor*›
⟨*proof*⟩
lift-definition *mask-uint64* :: ‹*nat* ⇒ *uint64*› **is** *mask* ⟨*proof*⟩
lift-definition *push-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *push-bit* ⟨*proof*⟩
lift-definition *drop-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *drop-bit* ⟨*proof*⟩
lift-definition *signed-drop-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *signed-drop-bit*
⟨*proof*⟩
lift-definition *take-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *take-bit* ⟨*proof*⟩
lift-definition *set-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *Bit-Operations.set-bit*
⟨*proof*⟩
lift-definition *unset-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *unset-bit* ⟨*proof*⟩
lift-definition *flip-bit-uint64* :: ‹*nat* ⇒ *uint64* ⇒ *uint64*› **is** *flip-bit* ⟨*proof*⟩

global-interpretation *uint64*: *word-type-copy-bits Abs-uint64 Rep-uint64 signed-drop-bit-uint64*
⟨*proof*⟩

instance
⟨*proof*⟩

end

lift-definition *uint64-of-nat* :: ‹*nat* ⇒ *uint64*›
is *word-of-nat* ⟨*proof*⟩

lift-definition *nat-of-uint64* :: ‹*uint64* ⇒ *nat*›
is *unat* ⟨*proof*⟩

lift-definition *uint64-of-int* :: ‹*int* ⇒ *uint64*›
is *word-of-int* ⟨*proof*⟩

lift-definition *int-of-uint64* :: ‹*uint64* ⇒ *int*›
is *uint* ⟨*proof*⟩

context
includes *integer.lifting*
begin

lift-definition *Uint64* :: ‹*integer* ⇒ *uint64*›
is *word-of-int* ⟨*proof*⟩

lift-definition *integer-of-uint64* :: ‹*uint64* ⇒ *integer*›
is *uint* ⟨*proof*⟩

end

global-interpretation *uint64*: *word-type-copy-more Abs-uint64 Rep-uint64 signed-drop-bit-uint64 uint64-of-nat nat-of-uint64 uint64-of-int int-of-uint64 Uint64 integer-of-uint64*
⟨*proof*⟩

instantiation *uint64* :: {*size*, *msb*, *lsb*, *set-bit*, *bit-comprehension*}
begin

lift-definition *size-uint64* :: ‹*uint64* ⇒ *nat*› **is** *size* ⟨*proof*⟩

lift-definition *msb-uint64* :: ‹*uint64* ⇒ *bool*› **is** *msb* ⟨*proof*⟩
lift-definition *lsb-uint64* :: ‹*uint64* ⇒ *bool*› **is** *lsb* ⟨*proof*⟩

Workaround: avoid name space clash by spelling out *lift-definition* explicitly.

definition *set-bit-uint64* :: ‹*uint64* ⇒ *nat* ⇒ *bool* ⇒ *uint64*›
where *set-bit-uint64-eq*: ‹*set-bit-uint64 a n b* = (*if b then Bit-Operations.set-bit else unset-bit*) *n a*›

context
includes *lifting-syntax*
begin

lemma *set-bit-uint64-transfer* [*transfer-rule*]:
‹(*cr-uint64* ===> (=) ===> (⟷) ===> *cr-uint64*) *Generic-set-bit.set-bit Generic-set-bit.set-bit*›
⟨*proof*⟩

end

lift-definition *set-bits-uint64* :: ‹(*nat* ⇒ *bool*) ⇒ *uint64*› **is** *set-bits* ⟨*proof*⟩
lift-definition *set-bits-aux-uint64* :: ‹(*nat* ⇒ *bool*) ⇒ *nat* ⇒ *uint64* ⇒ *uint64*› **is** *set-bits-aux* ⟨*proof*⟩

global-interpretation *uint64*: *word-type-copy-misc Abs-uint64 Rep-uint64 signed-drop-bit-uint64 uint64-of-nat nat-of-uint64 uint64-of-int int-of-uint64 Uint64 integer-of-uint64 64 set-bits-aux-uint64*
⟨*proof*⟩

instance ⟨*proof*⟩

end

6.2 Code setup

For SML, we generate an implementation of unsigned 64-bit words using `IntInf.int`. If `LargeWord.wordSize > 63` of the Isabelle/ML runtime environment holds, then we assume that there is also a *Word64* structure available and accordingly replace the implementation for the target `Eval`.

code-printing code-module *Uint64* ⇀ (*SML*) ‹(∗ *Test that words can handle numbers between 0 and 63* ∗)
val - = if 6 <= Word.wordSize then () else raise (Fail (wordSize less than 6));

structure Uint64 : sig
eqtype uint64;
val zero : uint64;
val one : uint64;
val fromInt : IntInf.int -> uint64;
val toInt : uint64 -> IntInf.int;
val toLarge : uint64 -> LargeWord.word;
val fromLarge : LargeWord.word -> uint64
val plus : uint64 -> uint64 -> uint64;
val minus : uint64 -> uint64 -> uint64;
val times : uint64 -> uint64 -> uint64;
val divide : uint64 -> uint64 -> uint64;
val modulus : uint64 -> uint64 -> uint64;
val negate : uint64 -> uint64;
val less-eq : uint64 -> uint64 -> bool;
val less : uint64 -> uint64 -> bool;
val notb : uint64 -> uint64;
val andb : uint64 -> uint64 -> uint64;
val orb : uint64 -> uint64 -> uint64;
val xorb : uint64 -> uint64 -> uint64;
val shiftl : uint64 -> IntInf.int -> uint64;
val shiftr : uint64 -> IntInf.int -> uint64;
val shiftr-signed : uint64 -> IntInf.int -> uint64;
val set-bit : uint64 -> IntInf.int -> bool -> uint64;
val test-bit : uint64 -> IntInf.int -> bool;

```
end = struct

type uint64 = IntInf.int;

val mask = 0xFFFFFFFFFFFFFFFF : IntInf.int;

val zero = 0 : IntInf.int;

val one = 1 : IntInf.int;

fun fromInt x = IntInf.andb(x, mask);

fun toInt x = x

fun toLarge x = LargeWord.fromLargeInt (IntInf.toLarge x);

fun fromLarge x = IntInf.fromLarge (LargeWord.toLargeInt x);

fun plus x y = IntInf.andb(IntInf.+(x, y), mask);

fun minus x y = IntInf.andb(IntInf.-(x, y), mask);

fun negate x = IntInf.andb(IntInf.~(x), mask);

fun times x y = IntInf.andb(IntInf.*(x, y), mask);

fun divide x y = IntInf.div(x, y);

fun modulus x y = IntInf.mod(x, y);

fun less-eq x y = IntInf.<=(x, y);

fun less x y = IntInf.<(x, y);

fun notb x = IntInf.andb(IntInf.notb(x), mask);

fun orb x y = IntInf.orb(x, y);

fun andb x y = IntInf.andb(x, y);

fun xorb x y = IntInf.xorb(x, y);

val maxWord = IntInf.pow (2, Word.wordSize);

fun shiftl x n =
  if n < maxWord then IntInf.andb(IntInf.<< (x, Word.fromLargeInt (IntInf.toLarge
n)), mask)
   else 0;
```

```
fun shiftr x n =
  if n < maxWord then IntInf.~>> (x, Word.fromLargeInt (IntInf.toLarge n))
  else 0;

val msb-mask = 0x8000000000000000 : IntInf.int;

fun shiftr-signed x i =
  if IntInf.andb(x, msb-mask) = 0 then shiftr x i
  else if i >= 64 then 0xFFFFFFFFFFFFFFFF
  else let
    val x' = shiftr x i
    val m' = IntInf.andb(IntInf.<<(mask, Word.max(0w64 - Word.fromLargeInt
(IntInf.toLarge i), 0w0)), mask)
  in IntInf.orb(x', m') end;

fun test-bit x n =
 if n < maxWord then IntInf.andb (x, IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge
n))) <> 0
  else false;

fun set-bit x n b =
  if n < 64 then
    if b then IntInf.orb (x, IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge n)))
   else IntInf.andb (x, IntInf.notb (IntInf.<< (1, Word.fromLargeInt (IntInf.toLarge
n))))
  else x;

end
›
code-reserved SML Uint64

⟨ML⟩

code-printing code-module Uint64 → (Haskell)
 ‹module Uint64(Int64, Word64) where

  import Data.Int(Int64)
  import Data.Word(Word64)›
code-reserved Haskell Uint64
```

OCaml and Scala provide only signed 64bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

```
code-printing code-module Uint64 → (OCaml)
‹module Uint64 : sig
  val less : int64 -> int64 -> bool
  val less-eq : int64 -> int64 -> bool
  val set-bit : int64 -> Z.t -> bool -> int64
  val shiftl : int64 -> Z.t -> int64
  val shiftr : int64 -> Z.t -> int64
```

```
  val shiftr-signed : int64 -> Z.t -> int64
  val test-bit : int64 -> Z.t -> bool
end = struct

(* negative numbers have their highest bit set,
   so they are greater than positive ones *)
let less x y =
  if Int64.compare x Int64.zero < 0 then
    Int64.compare y Int64.zero < 0 && Int64.compare x y < 0
  else Int64.compare y Int64.zero < 0 || Int64.compare x y < 0;;

let less-eq x y =
  if Int64.compare x Int64.zero < 0 then
    Int64.compare y Int64.zero < 0 && Int64.compare x y <= 0
  else Int64.compare y Int64.zero < 0 || Int64.compare x y <= 0;;

let set-bit x n b =
  let mask = Int64.shift-left Int64.one (Z.to-int n)
  in if b then Int64.logor x mask
     else Int64.logand x (Int64.lognot mask);;

let shiftl x n = Int64.shift-left x (Z.to-int n);;

let shiftr x n = Int64.shift-right-logical x (Z.to-int n);;

let shiftr-signed x n = Int64.shift-right x (Z.to-int n);;

let test-bit x n =
  Int64.compare
    (Int64.logand x (Int64.shift-left Int64.one (Z.to-int n)))
    Int64.zero
  <> 0;;

end;; (*struct Uint64*)›
code-reserved OCaml Uint64

code-printing code-module Uint64 → (Scala)
‹object Uint64 {

def less(x: Long, y: Long) : Boolean =
  if (x < 0) y < 0 && x < y
  else y < 0 || x < y

def less-eq(x: Long, y: Long) : Boolean =
  if (x < 0) y < 0 && x <= y
  else y < 0 || x <= y

def set-bit(x: Long, n: BigInt, b: Boolean) : Long =
  if (b)
```

```
    x | (1L << n.intValue)
  else
    x & (1L << n.intValue).unary-~

def shiftl(x: Long, n: BigInt) : Long = x << n.intValue

def shiftr(x: Long, n: BigInt) : Long = x >>> n.intValue

def shiftr-signed(x: Long, n: BigInt) : Long = x >> n.intValue

def test-bit(x: Long, n: BigInt) : Boolean =
  (x & (1L << n.intValue)) != 0

} /* object Uint64 */›
code-reserved Scala Uint64
```

OCaml's conversion from Big_int to int64 demands that the value fits int a signed 64-bit integer. The following justifies the implementation.

```
context
  includes bit-operations-syntax
begin

definition Uint64-signed :: integer ⇒ uint64
where Uint64-signed i = (if i < −(0x8000000000000000) ∨ i ≥ 0x8000000000000000
then undefined Uint64 i else Uint64 i)

lemma Uint64-code [code]:
  Uint64 i =
  (let i′ = i AND 0xFFFFFFFFFFFFFFFF
   in if bit i′ 63 then Uint64-signed (i′ − 0x10000000000000000) else Uint64-signed
i′)
  including undefined-transfer integer.lifting ⟨proof⟩

lemma Uint64-signed-code [code]:
  Rep-uint64 (Uint64-signed i) =
   (if i < −(0x8000000000000000) ∨ i ≥ 0x8000000000000000 then Rep-uint64
(undefined Uint64 i) else word-of-int (int-of-integer-symbolic i))
⟨proof⟩

end
```

Avoid *Abs-uint64* in generated code, use *Rep-uint64′* instead. The symbolic implementations for code_simp use *Rep-uint64*.

The new destructor *Rep-uint64′* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint64* directly, because that makes code_simp loop.

If code generation raises Match, some equation probably contains *Rep-uint64* ([code abstract] equations for *uint64* may use *Rep-uint64* because these

instances will be folded away.)

To convert *64 word* values into *uint64*, use *Abs-uint64'*.

definition *Rep-uint64'* **where** [*simp*]: *Rep-uint64' = Rep-uint64*

lemma *Rep-uint64'-transfer* [*transfer-rule*]:
rel-fun cr-uint64 (=) (λx. x) Rep-uint64'
⟨*proof*⟩

lemma *Rep-uint64'-code* [*code*]: *Rep-uint64' x = (BITS n. bit x n)*
⟨*proof*⟩

lift-definition *Abs-uint64' :: 64 word ⇒ uint64* **is** *λx :: 64 word. x* ⟨*proof*⟩

lemma *Abs-uint64'-code* [*code*]:
Abs-uint64' x = Uint64 (integer-of-int (uint x))
including *integer.lifting* ⟨*proof*⟩

declare [[*code drop*: *term-of-class.term-of :: uint64 ⇒* -]]

lemma *term-of-uint64-code* [*code*]:
defines *TR ≡ typerep.Typerep* **and** *bit0 ≡ STR ''Numeral-Type.bit0''*
shows
term-of-class.term-of x =
Code-Evaluation.App (Code-Evaluation.Const (STR ''Uint64.uint64.Abs-uint64'')
(TR (STR ''fun'') [TR (STR ''Word.word'') [TR bit0 [TR bit0 [TR bit0 [TR bit0
[TR bit0 [TR bit0 [TR (STR ''Numeral-Type.num1'') []]]]]]]], TR (STR ''Uint64.uint64'')
[]]))
(term-of-class.term-of (Rep-uint64' x))
⟨*proof*⟩

code-printing
type-constructor *uint64* ⇀
(*SML*) *Uint64.uint64* **and**
(*Haskell*) *Uint64.Word64* **and**
(*OCaml*) *int64* **and**
(*Scala*) *Long*
| **constant** *Uint64* ⇀
(*SML*) *Uint64.fromInt* **and**
(*Haskell*) (*Prelude.fromInteger* - :: *Uint64.Word64*) **and**
(*Haskell-Quickcheck*) (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint64.Word64*)
and
(*Scala*) -.*longValue*
| **constant** *Uint64-signed* ⇀
(*OCaml*) *Z.to'-int64*
| **constant** *0 :: uint64* ⇀
(*SML*) *Uint64.zero* **and**
(*Haskell*) (*0 :: Uint64.Word64*) **and**
(*OCaml*) *Int64.zero* **and**
(*Scala*) *0*

```
| constant 1 :: uint64 ⇀
  (SML) Uint64.one and
  (Haskell) (1 :: Uint64.Word64) and
  (OCaml) Int64.one and
  (Scala) 1
| constant plus :: uint64 ⇒ - ⇀
  (SML) Uint64.plus and
  (Haskell) infixl 6 + and
  (OCaml) Int64.add and
  (Scala) infixl 7 +
| constant uminus :: uint64 ⇒ - ⇀
  (SML) Uint64.negate and
  (Haskell) negate and
  (OCaml) Int64.neg and
  (Scala) !(− -)
| constant minus :: uint64 ⇒ - ⇀
  (SML) Uint64.minus and
  (Haskell) infixl 6 − and
  (OCaml) Int64.sub and
  (Scala) infixl 7 −
| constant times :: uint64 ⇒ - ⇒ - ⇀
  (SML) Uint64.times and
  (Haskell) infixl 7 ∗ and
  (OCaml) Int64.mul and
  (Scala) infixl 8 ∗
| constant HOL.equal :: uint64 ⇒ - ⇒ bool ⇀
  (SML) !((- : Uint64.uint64) = -) and
  (Haskell) infix 4 == and
  (OCaml) (Int64.compare - - = 0) and
  (Scala) infixl 5 ==
| class-instance uint64 :: equal ⇀
  (Haskell) −
| constant less-eq :: uint64 ⇒ - ⇒ bool ⇀
  (SML) Uint64.less'-eq and
  (Haskell) infix 4 <= and
  (OCaml) Uint64.less'-eq and
  (Scala) Uint64.less'-eq
| constant less :: uint64 ⇒ - ⇒ bool ⇀
  (SML) Uint64.less and
  (Haskell) infix 4 < and
  (OCaml) Uint64.less and
  (Scala) Uint64.less
| constant Bit-Operations.not :: uint64 ⇒ - ⇀
  (SML) Uint64.notb and
  (Haskell) Data'-Bits.complement and
  (OCaml) Int64.lognot and
  (Scala) -.unary'-~
| constant Bit-Operations.and :: uint64 ⇒ - ⇀
  (SML) Uint64.andb and
```

```
  (Haskell) infixl 7 Data-Bits..&. and
  (OCaml) Int64.logand and
  (Scala) infixl 3 &
| constant Bit-Operations.or :: uint64 ⇒ - ⇀
  (SML) Uint64.orb and
  (Haskell) infixl 5 Data-Bits..|. and
  (OCaml) Int64.logor and
  (Scala) infixl 1 |
| constant Bit-Operations.xor :: uint64 ⇒ - ⇀
  (SML) Uint64.xorb and
  (Haskell) Data'-Bits.xor and
  (OCaml) Int64.logxor and
  (Scala) infixl 2 ^

definition uint64-divmod :: uint64 ⇒ uint64 ⇒ uint64 × uint64 where
  uint64-divmod x y =
  (if y = 0 then (undefined ((div) :: uint64 ⇒ -) x (0 :: uint64), undefined ((mod)
:: uint64 ⇒ -) x (0 :: uint64))
  else (x div y, x mod y))

definition uint64-div :: uint64 ⇒ uint64 ⇒ uint64
where uint64-div x y = fst (uint64-divmod x y)

definition uint64-mod :: uint64 ⇒ uint64 ⇒ uint64
where uint64-mod x y = snd (uint64-divmod x y)

lemma div-uint64-code [code]: x div y = (if y = 0 then 0 else uint64-div x y)
including undefined-transfer ⟨proof⟩

lemma mod-uint64-code [code]: x mod y = (if y = 0 then x else uint64-mod x y)
including undefined-transfer ⟨proof⟩

definition uint64-sdiv :: uint64 ⇒ uint64 ⇒ uint64
where [code del]:
  uint64-sdiv x y =
   (if y = 0 then undefined ((div) :: uint64 ⇒ -) x (0 :: uint64)
    else Abs-uint64 (Rep-uint64 x sdiv Rep-uint64 y))

definition div0-uint64 :: uint64 ⇒ uint64
where [code del]: div0-uint64 x = undefined ((div) :: uint64 ⇒ -) x (0 :: uint64)
declare [[code abort: div0-uint64]]

definition mod0-uint64 :: uint64 ⇒ uint64
where [code del]: mod0-uint64 x = undefined ((mod) :: uint64 ⇒ -) x (0 :: uint64)
declare [[code abort: mod0-uint64]]

lemma uint64-divmod-code [code]:
  uint64-divmod x y =
  (if 0x8000000000000000 ≤ y then if x < y then (0, x) else (1, x − y)
```

```
  else if y = 0 then (div0-uint64 x, mod0-uint64 x)
  else let q = push-bit 1 (uint64-sdiv (drop-bit 1 x) y);
           r = x - q * y
       in if r ≥ y then (q + 1, r - y) else (q, r))
 including undefined-transfer ⟨proof⟩

lemma uint64-sdiv-code [code]:
  Rep-uint64 (uint64-sdiv x y) =
   (if y = 0 then Rep-uint64 (undefined ((div) :: uint64 ⇒ -) x (0 :: uint64))
    else Rep-uint64 x sdiv Rep-uint64 y)
⟨proof⟩
```

Note that we only need a translation for signed division, but not for the remainder because *uint64-divmod ?x ?y = (if 9223372036854775808 ≤ ?y then if ?x < ?y then (0, ?x) else (1, ?x − ?y) else if ?y = 0 then (div0-uint64 ?x, mod0-uint64 ?x) else let q = push-bit 1 (uint64-sdiv (drop-bit 1 ?x) ?y); r = ?x − q * ?y in if ?y ≤ r then (q + 1, r − ?y) else (q, r))* computes both with division only.

```
code-printing
  constant uint64-div ⇀
  (SML) Uint64.divide and
  (Haskell) Prelude.div
| constant uint64-mod ⇀
  (SML) Uint64.modulus and
  (Haskell) Prelude.mod
| constant uint64-divmod ⇀
  (Haskell) divmod
| constant uint64-sdiv ⇀
  (OCaml) Int64.div and
  (Scala) - '/ -

definition uint64-test-bit :: uint64 ⇒ integer ⇒ bool
where [code del]:
  uint64-test-bit x n =
  (if n < 0 ∨ 63 < n then undefined (bit :: uint64 ⇒ -) x n
   else bit x (nat-of-integer n))

lemma bit-uint64-code [code]:
  bit x n ⟷ n < 64 ∧ uint64-test-bit x (integer-of-nat n)
  including undefined-transfer integer.lifting ⟨proof⟩

lemma uint64-test-bit-code [code]:
  uint64-test-bit w n =
  (if n < 0 ∨ 63 < n then undefined (bit :: uint64 ⇒ -) w n else bit (Rep-uint64
w) (nat-of-integer n))
  ⟨proof⟩

code-printing constant uint64-test-bit ⇀
  (SML) Uint64.test'-bit and
```

(*Haskell*) *Data′-Bits.testBitBounded* **and**
(*OCaml*) *Uint64.test′-bit* **and**
(*Scala*) *Uint64.test′-bit* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 64 then raise* (*Fail argument to uint64′-test′-bit out of bounds*) *else Uint64.test′-bit x i*)

definition *uint64-set-bit* :: *uint64* ⇒ *integer* ⇒ *bool* ⇒ *uint64*
where [*code del*]:
uint64-set-bit x n b =
(*if n < 0 ∨ 63 < n then undefined* (*set-bit* :: *uint64* ⇒ -) *x n b*
else set-bit x (*nat-of-integer n*) *b*)

lemma *set-bit-uint64-code* [*code*]:
set-bit x n b = (*if n < 64 then uint64-set-bit x* (*integer-of-nat n*) *b else x*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint64-set-bit-code* [*code*]:
Rep-uint64 (*uint64-set-bit w n b*) =
(*if n < 0 ∨ 63 < n then Rep-uint64* (*undefined* (*set-bit* :: *uint64* ⇒ -) *w n b*)
else set-bit (*Rep-uint64 w*) (*nat-of-integer n*) *b*)
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint64-set-bit* ⇀
(*SML*) *Uint64.set′-bit* **and**
(*Haskell*) *Data′-Bits.setBitBounded* **and**
(*OCaml*) *Uint64.set′-bit* **and**
(*Scala*) *Uint64.set′-bit* **and**
(*Eval*) (*fn x => fn i => fn b => if i < 0 orelse i >= 64 then raise* (*Fail argument to uint64′-set′-bit out of bounds*) *else Uint64.set′-bit x i b*)

definition *uint64-shiftl* :: *uint64* ⇒ *integer* ⇒ *uint64*
where [*code del*]:
uint64-shiftl x n = (*if n < 0 ∨ 64 ≤ n then undefined* (*push-bit* :: *nat* ⇒ *uint64* ⇒ -) *x n else push-bit* (*nat-of-integer n*) *x*)

lemma *shiftl-uint64-code* [*code*]: *push-bit n x* = (*if n < 64 then uint64-shiftl x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint64-shiftl-code* [*code*]:
Rep-uint64 (*uint64-shiftl w n*) =
(*if n < 0 ∨ 64 ≤ n then Rep-uint64* (*undefined* (*push-bit* :: *nat* ⇒ *uint64* ⇒ -) *w n*) *else push-bit* (*nat-of-integer n*) (*Rep-uint64 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint64-shiftl* ⇀
(*SML*) *Uint64.shiftl* **and**
(*Haskell*) *Data′-Bits.shiftlBounded* **and**
(*OCaml*) *Uint64.shiftl* **and**

(*Scala*) *Uint64.shiftl* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 64 then raise* (*Fail argument to uint64'-shiftl out of bounds*) *else Uint64.shiftl x i*)

definition *uint64-shiftr* :: *uint64* ⇒ *integer* ⇒ *uint64*
where [*code del*]:
uint64-shiftr x n = (*if n < 0* ∨ *64* ≤ *n then undefined* (*drop-bit* :: *nat* ⇒ *uint64* ⇒ -) *x n else drop-bit* (*nat-of-integer n*) *x*)

lemma *shiftr-uint64-code* [*code*]: *drop-bit n x* = (*if n < 64 then uint64-shiftr x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint64-shiftr-code* [*code*]:
Rep-uint64 (*uint64-shiftr w n*) =
(*if n < 0* ∨ *64* ≤ *n then Rep-uint64* (*undefined* (*drop-bit* :: *nat* ⇒ *uint64* ⇒ -) *w n*) *else drop-bit* (*nat-of-integer n*) (*Rep-uint64 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint64-shiftr* ⇀
(*SML*) *Uint64.shiftr* **and**
(*Haskell*) *Data'-Bits.shiftrBounded* **and**
(*OCaml*) *Uint64.shiftr* **and**
(*Scala*) *Uint64.shiftr* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 64 then raise* (*Fail argument to uint64'-shiftr out of bounds*) *else Uint64.shiftr x i*)

definition *uint64-sshiftr* :: *uint64* ⇒ *integer* ⇒ *uint64*
where [*code del*]:
uint64-sshiftr x n =
(*if n < 0* ∨ *64* ≤ *n then undefined signed-drop-bit-uint64 n x else signed-drop-bit-uint64* (*nat-of-integer n*) *x*)

lemma *sshiftr-uint64-code* [*code*]:
signed-drop-bit-uint64 n x =
(*if n < 64 then uint64-sshiftr x* (*integer-of-nat n*) *else if bit x 63 then* − *1 else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint64-sshiftr-code* [*code*]:
Rep-uint64 (*uint64-sshiftr w n*) =
(*if n < 0* ∨ *64* ≤ *n then Rep-uint64* (*undefined signed-drop-bit-uint64 n w*) *else signed-drop-bit* (*nat-of-integer n*) (*Rep-uint64 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint64-sshiftr* ⇀
(*SML*) *Uint64.shiftr'-signed* **and**
(*Haskell*)
(*Prelude.fromInteger* (*Prelude.toInteger* (*Data'-Bits.shiftrBounded* (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint64.Int64*) -)) :: *Uint64.Word64*) **and**

(*OCaml*) *Uint64.shiftr'-signed* **and**
(*Scala*) *Uint64.shiftr'-signed* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 64 then raise* (*Fail argument to uint64'-shiftr'-signed out of bounds*) *else Uint64.shiftr'-signed x i*)

context
includes *bit-operations-syntax*
begin

lemma *uint64-msb-test-bit*: *msb x* ⟷ *bit* (*x* :: *uint64*) *63*
⟨*proof*⟩

lemma *msb-uint64-code* [*code*]: *msb x* ⟷ *uint64-test-bit x 63*
⟨*proof*⟩

lemma *uint64-of-int-code* [*code*]:
uint64-of-int i = *Uint64* (*integer-of-int i*)
including *integer.lifting* ⟨*proof*⟩

lemma *int-of-uint64-code* [*code*]:
int-of-uint64 x = *int-of-integer* (*integer-of-uint64 x*)
including *integer.lifting* ⟨*proof*⟩

lemma *uint64-of-nat-code* [*code*]:
uint64-of-nat = *uint64-of-int* ∘ *int*
⟨*proof*⟩

lemma *nat-of-uint64-code* [*code*]:
nat-of-uint64 x = *nat-of-integer* (*integer-of-uint64 x*)
⟨*proof*⟩ **including** *integer.lifting* ⟨*proof*⟩

definition *integer-of-uint64-signed* :: *uint64* ⇒ *integer*
where
integer-of-uint64-signed n = (*if bit n 63 then undefined integer-of-uint64 n else integer-of-uint64 n*)

lemma *integer-of-uint64-signed-code* [*code*]:
integer-of-uint64-signed n =
(*if bit n 63 then undefined integer-of-uint64 n else integer-of-int* (*uint* (*Rep-uint64' n*)))
⟨*proof*⟩

lemma *integer-of-uint64-code* [*code*]:
integer-of-uint64 n =
(*if bit n 63 then integer-of-uint64-signed* (*n AND 0x7FFFFFFFFFFFFFFF*) *OR 0x8000000000000000 else integer-of-uint64-signed n*)
⟨*proof*⟩
including *integer.lifting* ⟨*proof*⟩

end

code-printing
  constant *integer-of-uint64* $\rightharpoonup$
  (*SML*) *Uint64.toInt* **and**
  (*Haskell*) *Prelude.toInteger*
| **constant** *integer-of-uint64-signed* $\rightharpoonup$
  (*OCaml*) *Z.of'-int64* **and**
  (*Scala*) *BigInt*

6.3 Quickcheck setup

definition *uint64-of-natural* :: *natural* $\Rightarrow$ *uint64*
where *uint64-of-natural x* $\equiv$ *Uint64* (*integer-of-natural x*)

instantiation *uint64* :: {*random*, *exhaustive*, *full-exhaustive*} **begin**
definition *random-uint64* $\equiv$ *qc-random-cnv uint64-of-natural*
definition *exhaustive-uint64* $\equiv$ *qc-exhaustive-cnv uint64-of-natural*
definition *full-exhaustive-uint64* $\equiv$ *qc-full-exhaustive-cnv uint64-of-natural*
instance ⟨*proof*⟩
end

instantiation *uint64* :: *narrowing* **begin**

interpretation *quickcheck-narrowing-samples*
  $\lambda i.$ *let x* = *Uint64 i in* (*x*, *0xFFFFFFFFFFFFFFFF* − *x*) *0*
  Typerep.Typerep (*STR ''Uint64.uint64''*) [] ⟨*proof*⟩

definition *narrowing-uint64 d* = *qc-narrowing-drawn-from* (*narrowing-samples d*) *d*
declare [[*code drop*: *partial-term-of* :: *uint64 itself* $\Rightarrow$ -]]
lemmas *partial-term-of-uint64* [*code*] = *partial-term-of-code*

instance ⟨*proof*⟩
end

end

Chapter 7

Unsigned words of 32 bits

theory *Uint32* **imports**
Code-Target-Word-Base Word-Type-Copies
begin

7.1 Type definition and primitive operations

typedef *uint32* = ‹*UNIV* :: *32 word set*› ⟨*proof*⟩

global-interpretation *uint32*: *word-type-copy Abs-uint32 Rep-uint32*
⟨*proof*⟩

setup-lifting *type-definition-uint32*

declare *uint32.of-word-of* [*code abstype*]

declare *Quotient-uint32* [*transfer-rule*]

instantiation *uint32* :: ‹{*comm-ring-1*, *semiring-modulo*, *equal*, *linorder*}›
begin

lift-definition *zero-uint32* :: *uint32* **is** *0* ⟨*proof*⟩
lift-definition *one-uint32* :: *uint32* **is** *1* ⟨*proof*⟩
lift-definition *plus-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹(+)› ⟨*proof*⟩
lift-definition *uminus-uint32* :: ‹*uint32* ⇒ *uint32*› **is** *uminus* ⟨*proof*⟩
lift-definition *minus-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹(−)› ⟨*proof*⟩
lift-definition *times-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹(∗)› ⟨*proof*⟩
lift-definition *divide-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹(*div*)› ⟨*proof*⟩
lift-definition *modulo-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹(*mod*)› ⟨*proof*⟩
lift-definition *equal-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *bool*› **is** ‹*HOL.equal*› ⟨*proof*⟩
lift-definition *less-eq-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *bool*› **is** ‹(≤)› ⟨*proof*⟩
lift-definition *less-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *bool*› **is** ‹(<)› ⟨*proof*⟩

global-interpretation *uint32*: *word-type-copy-ring Abs-uint32 Rep-uint32*
⟨*proof*⟩

instance ⟨*proof*⟩

end

instantiation *uint32* :: *ring-bit-operations*
begin

lift-definition *bit-uint32* :: ‹*uint32* ⇒ *nat* ⇒ *bool*› **is** *bit* ⟨*proof*⟩
lift-definition *not-uint32* :: ‹*uint32* ⇒ *uint32*› **is** ‹*Bit-Operations.not*› ⟨*proof*⟩
lift-definition *and-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹*Bit-Operations.and*›
⟨*proof*⟩
lift-definition *or-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹*Bit-Operations.or*›
⟨*proof*⟩
lift-definition *xor-uint32* :: ‹*uint32* ⇒ *uint32* ⇒ *uint32*› **is** ‹*Bit-Operations.xor*›
⟨*proof*⟩
lift-definition *mask-uint32* :: ‹*nat* ⇒ *uint32*› **is** *mask* ⟨*proof*⟩
lift-definition *push-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *push-bit* ⟨*proof*⟩
lift-definition *drop-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *drop-bit* ⟨*proof*⟩
lift-definition *signed-drop-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *signed-drop-bit*
⟨*proof*⟩
lift-definition *take-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *take-bit* ⟨*proof*⟩
lift-definition *set-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *Bit-Operations.set-bit*
⟨*proof*⟩
lift-definition *unset-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *unset-bit* ⟨*proof*⟩
lift-definition *flip-bit-uint32* :: ‹*nat* ⇒ *uint32* ⇒ *uint32*› **is** *flip-bit* ⟨*proof*⟩

global-interpretation *uint32*: *word-type-copy-bits Abs-uint32 Rep-uint32 signed-drop-bit-uint32*
⟨*proof*⟩

instance
⟨*proof*⟩

end

lift-definition *uint32-of-nat* :: ‹*nat* ⇒ *uint32*›
is *word-of-nat* ⟨*proof*⟩

lift-definition *nat-of-uint32* :: ‹*uint32* ⇒ *nat*›
is *unat* ⟨*proof*⟩

lift-definition *uint32-of-int* :: ‹*int* ⇒ *uint32*›
is *word-of-int* ⟨*proof*⟩

lift-definition *int-of-uint32* :: ‹*uint32* ⇒ *int*›
is *uint* ⟨*proof*⟩

context
includes *integer.lifting*

begin

lift-definition *Uint32* :: ‹*integer* ⇒ *uint32*›
  is *word-of-int* ⟨*proof*⟩

lift-definition *integer-of-uint32* :: ‹*uint32* ⇒ *integer*›
  is *uint* ⟨*proof*⟩

end

global-interpretation *uint32*: *word-type-copy-more Abs-uint32 Rep-uint32 signed-drop-bit-uint32*
  uint32-of-nat nat-of-uint32 uint32-of-int int-of-uint32 Uint32 integer-of-uint32
  ⟨*proof*⟩

instantiation *uint32* :: {*size*, *msb*, *lsb*, *set-bit*, *bit-comprehension*}
begin

lift-definition *size-uint32* :: ‹*uint32* ⇒ *nat*› **is** *size* ⟨*proof*⟩

lift-definition *msb-uint32* :: ‹*uint32* ⇒ *bool*› **is** *msb* ⟨*proof*⟩
lift-definition *lsb-uint32* :: ‹*uint32* ⇒ *bool*› **is** *lsb* ⟨*proof*⟩

Workaround: avoid name space clash by spelling out *lift-definition* explicitly.

definition *set-bit-uint32* :: ‹*uint32* ⇒ *nat* ⇒ *bool* ⇒ *uint32*›
  where *set-bit-uint32-eq*: ‹*set-bit-uint32 a n b* = (*if b then Bit-Operations.set-bit*
else unset-bit) *n a*›

context
  includes *lifting-syntax*
begin

lemma *set-bit-uint32-transfer* [*transfer-rule*]:
  ‹(*cr-uint32* ===> (=) ===> (⟷) ===> *cr-uint32*) *Generic-set-bit.set-bit*
Generic-set-bit.set-bit›
  ⟨*proof*⟩

end

lift-definition *set-bits-uint32* :: ‹(*nat* ⇒ *bool*) ⇒ *uint32*› **is** *set-bits* ⟨*proof*⟩
lift-definition *set-bits-aux-uint32* :: ‹(*nat* ⇒ *bool*) ⇒ *nat* ⇒ *uint32* ⇒ *uint32*› **is**
set-bits-aux ⟨*proof*⟩

global-interpretation *uint32*: *word-type-copy-misc Abs-uint32 Rep-uint32 signed-drop-bit-uint32*
  uint32-of-nat nat-of-uint32 uint32-of-int int-of-uint32 Uint32 integer-of-uint32 32
set-bits-aux-uint32
  ⟨*proof*⟩

instance ⟨*proof*⟩

end

7.2 Code setup

code-printing code-module *Uint32* ⇀ (*SML*)
‹(∗ *Test that words can handle numbers between 0 and 31* ∗)
val - = if 5 <= Word.wordSize then () else raise (Fail (wordSize less than 5));

structure Uint32 : sig
val set-bit : Word32.word -> IntInf.int -> bool -> Word32.word
val shiftl : Word32.word -> IntInf.int -> Word32.word
val shiftr : Word32.word -> IntInf.int -> Word32.word
val shiftr-signed : Word32.word -> IntInf.int -> Word32.word
val test-bit : Word32.word -> IntInf.int -> bool
end = struct

fun set-bit x n b =
let val mask = Word32.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))
in if b then Word32.orb (x, mask)
else Word32.andb (x, Word32.notb mask)
end

fun shiftl x n =
Word32.<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
Word32.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
Word32.~>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
Word32.andb (x, Word32.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n)))
<> Word32.fromInt 0

end; (∗ *struct Uint32* ∗)›
code-reserved *SML Uint32*

code-printing code-module *Uint32* ⇀ (*Haskell*)
‹*module Uint32*(*Int32*, *Word32*) *where*

import Data.Int(*Int32*)
import Data.Word(*Word32*)›
code-reserved *Haskell Uint32*

OCaml and Scala provide only signed 32bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

code-printing code-module *Uint32* ⇀ (*OCaml*)
‹*module Uint32 : sig*

```
  val less : int32 -> int32 -> bool
  val less-eq : int32 -> int32 -> bool
  val set-bit : int32 -> Z.t -> bool -> int32
  val shiftl : int32 -> Z.t -> int32
  val shiftr : int32 -> Z.t -> int32
  val shiftr-signed : int32 -> Z.t -> int32
  val test-bit : int32 -> Z.t -> bool
end = struct

(* negative numbers have their highest bit set,
   so they are greater than positive ones *)
let less x y =
  if Int32.compare x Int32.zero < 0 then
    Int32.compare y Int32.zero < 0 && Int32.compare x y < 0
  else Int32.compare y Int32.zero < 0 || Int32.compare x y < 0;;

let less-eq x y =
  if Int32.compare x Int32.zero < 0 then
    Int32.compare y Int32.zero < 0 && Int32.compare x y <= 0
  else Int32.compare y Int32.zero < 0 || Int32.compare x y <= 0;;

let set-bit x n b =
  let mask = Int32.shift-left Int32.one (Z.to-int n)
  in if b then Int32.logor x mask
     else Int32.logand x (Int32.lognot mask);;

let shiftl x n = Int32.shift-left x (Z.to-int n);;

let shiftr x n = Int32.shift-right-logical x (Z.to-int n);;

let shiftr-signed x n = Int32.shift-right x (Z.to-int n);;

let test-bit x n =
  Int32.compare
    (Int32.logand x (Int32.shift-left Int32.one (Z.to-int n)))
    Int32.zero
  <> 0;;

end;; (*struct Uint32*)›
code-reserved OCaml Uint32

code-printing code-module Uint32 → (Scala)
‹object Uint32 {

def less(x: Int, y: Int) : Boolean =
  if (x < 0) y < 0 && x < y
  else y < 0 || x < y

def less-eq(x: Int, y: Int) : Boolean =
```

```
  if (x < 0) y < 0 && x <= y
  else y < 0 || x <= y

def set-bit(x: Int, n: BigInt, b: Boolean) : Int =
  if (b)
    x | (1 << n.intValue)
  else
    x & (1 << n.intValue).unary-~

def shiftl(x: Int, n: BigInt) : Int = x << n.intValue

def shiftr(x: Int, n: BigInt) : Int = x >>> n.intValue

def shiftr-signed(x: Int, n: BigInt) : Int = x >> n.intValue

def test-bit(x: Int, n: BigInt) : Boolean =
  (x & (1 << n.intValue)) != 0

} /* object Uint32 */›
code-reserved Scala Uint32
```

OCaml's conversion from Big_int to int32 demands that the value fits int a signed 32-bit integer. The following justifies the implementation.

```
context
  includes bit-operations-syntax
begin

definition Uint32-signed :: integer ⇒ uint32
where Uint32-signed i = (if i < −(0x80000000) ∨ i ≥ 0x80000000 then undefined
Uint32 i else Uint32 i)

lemma Uint32-code [code]:
  Uint32 i =
  (let i' = i AND 0xFFFFFFFF
   in if bit i' 31 then Uint32-signed (i' − 0x100000000) else Uint32-signed i')
  including undefined-transfer integer.lifting ⟨proof⟩

lemma Uint32-signed-code [code]:
  Rep-uint32 (Uint32-signed i) =
  (if i < −(0x80000000) ∨ i ≥ 0x80000000 then Rep-uint32 (undefined Uint32 i)
else word-of-int (int-of-integer-symbolic i))
⟨proof⟩

end
```

Avoid *Abs-uint32* in generated code, use *Rep-uint32'* instead. The symbolic implementations for code_simp use *Rep-uint32*.

The new destructor *Rep-uint32'* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint32* directly,

because that makes code_simp loop.

If code generation raises Match, some equation probably contains *Rep-uint32* ([code abstract] equations for *uint32* may use *Rep-uint32* because these instances will be folded away.)

To convert *32 word* values into *uint32*, use *Abs-uint32'*.

definition *Rep-uint32'* **where** [*simp*]: *Rep-uint32' = Rep-uint32*

lemma *Rep-uint32'-transfer* [*transfer-rule*]:
rel-fun cr-uint32 (=) (λx. x) Rep-uint32'
⟨*proof*⟩

lemma *Rep-uint32'-code* [*code*]: *Rep-uint32' x = (BITS n. bit x n)*
⟨*proof*⟩

lift-definition *Abs-uint32'* :: *32 word ⇒ uint32* **is** *λx* :: *32 word. x* ⟨*proof*⟩

lemma *Abs-uint32'-code* [*code*]:
Abs-uint32' x = Uint32 (integer-of-int (uint x))
including *integer.lifting* ⟨*proof*⟩

declare [[*code drop*: *term-of-class.term-of* :: *uint32 ⇒* -]]

lemma *term-of-uint32-code* [*code*]:
defines *TR ≡ typerep.Typerep* **and** *bit0 ≡ STR ''Numeral-Type.bit0''*
shows
term-of-class.term-of x =
Code-Evaluation.App (Code-Evaluation.Const (STR ''Uint32.uint32.Abs-uint32'')
(TR (STR ''fun'') [TR (STR ''Word.word'') [TR bit0 [TR bit0 [TR bit0 [TR bit0
[TR bit0 [TR (STR ''Numeral-Type.num1'') []]]]]]], TR (STR ''Uint32.uint32'')
[]]))
(term-of-class.term-of (Rep-uint32' x))
⟨*proof*⟩

code-printing
type-constructor *uint32* ⇀
(*SML*) *Word32.word* **and**
(*Haskell*) *Uint32.Word32* **and**
(*OCaml*) *int32* **and**
(*Scala*) *Int* **and**
(*Eval*) *Word32.word*
| **constant** *Uint32* ⇀
(*SML*) *Word32.fromLargeInt* (*IntInf.toLarge* -) **and**
(*Haskell*) (*Prelude.fromInteger* - :: *Uint32.Word32*) **and**
(*Haskell-Quickcheck*) (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint32.Word32*)
and
(*Scala*) -.*intValue*
| **constant** *Uint32-signed* ⇀
(*OCaml*) *Z.to'-int32*

```
| constant 0 :: uint32 ⇀
  (SML) (Word32.fromInt 0) and
  (Haskell) (0 :: Uint32.Word32) and
  (OCaml) Int32.zero and
  (Scala) 0
| constant 1 :: uint32 ⇀
  (SML) (Word32.fromInt 1) and
  (Haskell) (1 :: Uint32.Word32) and
  (OCaml) Int32.one and
  (Scala) 1
| constant plus :: uint32 ⇒ - ⇀
  (SML) Word32.+ ((-), (-)) and
  (Haskell) infixl 6 + and
  (OCaml) Int32.add and
  (Scala) infixl 7 +
| constant uminus :: uint32 ⇒ - ⇀
  (SML) Word32.~ and
  (Haskell) negate and
  (OCaml) Int32.neg and
  (Scala) !(− -)
| constant minus :: uint32 ⇒ - ⇀
  (SML) Word32.− ((-), (-)) and
  (Haskell) infixl 6 − and
  (OCaml) Int32.sub and
  (Scala) infixl 7 −
| constant times :: uint32 ⇒ - ⇒ - ⇀
  (SML) Word32.∗ ((-), (-)) and
  (Haskell) infixl 7 ∗ and
  (OCaml) Int32.mul and
  (Scala) infixl 8 ∗
| constant HOL.equal :: uint32 ⇒ - ⇒ bool ⇀
  (SML) !((- : Word32.word) = -) and
  (Haskell) infix 4 == and
  (OCaml) (Int32.compare - - = 0) and
  (Scala) infixl 5 ==
| class-instance uint32 :: equal ⇀
  (Haskell) −
| constant less-eq :: uint32 ⇒ - ⇒ bool ⇀
  (SML) Word32.<= ((-), (-)) and
  (Haskell) infix 4 <= and
  (OCaml) Uint32.less'-eq and
  (Scala) Uint32.less'-eq
| constant less :: uint32 ⇒ - ⇒ bool ⇀
  (SML) Word32.< ((-), (-)) and
  (Haskell) infix 4 < and
  (OCaml) Uint32.less and
  (Scala) Uint32.less
| constant Bit-Operations.not :: uint32 ⇒ - ⇀
  (SML) Word32.notb and
```

(*Haskell*) *Data′-Bits.complement* **and**
(*OCaml*) *Int32.lognot* **and**
(*Scala*) *-.unary′-*~
| **constant** *Bit-Operations.and* :: *uint32* ⇒ - ⇀
(*SML*) *Word32.andb* ((-),/ (-)) **and**
(*Haskell*) **infixl** *7 Data-Bits..&.* **and**
(*OCaml*) *Int32.logand* **and**
(*Scala*) **infixl** *3* &
| **constant** *Bit-Operations.or* :: *uint32* ⇒ - ⇀
(*SML*) *Word32.orb* ((-),/ (-)) **and**
(*Haskell*) **infixl** *5 Data-Bits..|.* **and**
(*OCaml*) *Int32.logor* **and**
(*Scala*) **infixl** *1* |
| **constant** *Bit-Operations.xor* :: *uint32* ⇒ - ⇀
(*SML*) *Word32.xorb* ((-),/ (-)) **and**
(*Haskell*) *Data′-Bits.xor* **and**
(*OCaml*) *Int32.logxor* **and**
(*Scala*) **infixl** *2* ^

definition *uint32-divmod* :: *uint32* ⇒ *uint32* ⇒ *uint32* × *uint32* **where**
uint32-divmod x y =
(*if y* = *0 then* (*undefined* ((*div*) :: *uint32* ⇒ -) *x* (*0* :: *uint32*), *undefined* ((*mod*) :: *uint32* ⇒ -) *x* (*0* :: *uint32*))
else (*x div y*, *x mod y*))

definition *uint32-div* :: *uint32* ⇒ *uint32* ⇒ *uint32*
where *uint32-div x y* = *fst* (*uint32-divmod x y*)

definition *uint32-mod* :: *uint32* ⇒ *uint32* ⇒ *uint32*
where *uint32-mod x y* = *snd* (*uint32-divmod x y*)

lemma *div-uint32-code* [*code*]: *x div y* = (*if y* = *0 then 0 else uint32-div x y*)
including *undefined-transfer* ⟨*proof*⟩

lemma *mod-uint32-code* [*code*]: *x mod y* = (*if y* = *0 then x else uint32-mod x y*)
including *undefined-transfer* ⟨*proof*⟩

definition *uint32-sdiv* :: *uint32* ⇒ *uint32* ⇒ *uint32*
where [*code del*]:
uint32-sdiv x y =
(*if y* = *0 then undefined* ((*div*) :: *uint32* ⇒ -) *x* (*0* :: *uint32*)
else Abs-uint32 (*Rep-uint32 x sdiv Rep-uint32 y*))

definition *div0-uint32* :: *uint32* ⇒ *uint32*
where [*code del*]: *div0-uint32 x* = *undefined* ((*div*) :: *uint32* ⇒ -) *x* (*0* :: *uint32*)
declare [[*code abort*: *div0-uint32*]]

definition *mod0-uint32* :: *uint32* ⇒ *uint32*
where [*code del*]: *mod0-uint32 x* = *undefined* ((*mod*) :: *uint32* ⇒ -) *x* (*0* :: *uint32*)

declare [[*code abort*: *mod0-uint32*]]

lemma *uint32-divmod-code* [*code*]:
uint32-divmod x y =
(*if 0x80000000* ≤ *y then if x* < *y then* (*0*, *x*) *else* (*1*, *x* − *y*)
else if y = *0 then* (*div0-uint32 x*, *mod0-uint32 x*)
else let q = *push-bit 1* (*uint32-sdiv* (*drop-bit 1 x*) *y*);
r = *x* − *q* ∗ *y*
in if r ≥ *y then* (*q* + *1*, *r* − *y*) *else* (*q*, *r*))
including *undefined-transfer* ⟨*proof*⟩

lemma *uint32-sdiv-code* [*code*]:
Rep-uint32 (*uint32-sdiv x y*) =
(*if y* = *0 then Rep-uint32* (*undefined* ((*div*) :: *uint32* ⇒ -) *x* (*0* :: *uint32*))
else Rep-uint32 x sdiv Rep-uint32 y)
⟨*proof*⟩

Note that we only need a translation for signed division, but not for the remainder because *uint32-divmod ?x ?y* = (*if 2147483648* ≤ *?y then if ?x* < *?y then* (*0*, *?x*) *else* (*1*, *?x* − *?y*) *else if ?y* = *0 then* (*div0-uint32 ?x*, *mod0-uint32 ?x*) *else let q* = *push-bit 1* (*uint32-sdiv* (*drop-bit 1 ?x*) *?y*); *r* = *?x* − *q* ∗ *?y in if ?y* ≤ *r then* (*q* + *1*, *r* − *?y*) *else* (*q*, *r*)) computes both with division only.

code-printing
constant *uint32-div* ⇀
(*SML*) *Word32.div* ((-), (-)) **and**
(*Haskell*) *Prelude.div*
| **constant** *uint32-mod* ⇀
(*SML*) *Word32.mod* ((-), (-)) **and**
(*Haskell*) *Prelude.mod*
| **constant** *uint32-divmod* ⇀
(*Haskell*) *divmod*
| **constant** *uint32-sdiv* ⇀
(*OCaml*) *Int32.div* **and**
(*Scala*) - *'/* -

definition *uint32-test-bit* :: *uint32* ⇒ *integer* ⇒ *bool*
where [*code del*]:
uint32-test-bit x n =
(*if n* < *0* ∨ *31* < *n then undefined* (*bit* :: *uint32* ⇒ -) *x n*
else bit x (*nat-of-integer n*))

lemma *test-bit-uint32-code* [*code*]:
bit x n ⟷ *n* < *32* ∧ *uint32-test-bit x* (*integer-of-nat n*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint32-test-bit-code* [*code*]:
uint32-test-bit w n =
(*if n* < *0* ∨ *31* < *n then undefined* (*bit* :: *uint32* ⇒ -) *w n else bit* (*Rep-uint32*

w) (*nat-of-integer n*))
⟨*proof*⟩

code-printing constant *uint32-test-bit* ⇀
(*SML*) *Uint32.test'-bit* **and**
(*Haskell*) *Data'-Bits.testBitBounded* **and**
(*OCaml*) *Uint32.test'-bit* **and**
(*Scala*) *Uint32.test'-bit* **and**
(*Eval*) (*fn w => fn n => if n < 0 orelse 32 <= n then raise* (*Fail argument to uint32'-test'-bit out of bounds*) *else Uint32.test'-bit w n*)

definition *uint32-set-bit* :: *uint32* ⇒ *integer* ⇒ *bool* ⇒ *uint32*
where [*code del*]:
uint32-set-bit x n b =
(*if n < 0* ∨ *31 < n then undefined* (*set-bit* :: *uint32* ⇒ -) *x n b*
else set-bit x (*nat-of-integer n*) *b*)

lemma *set-bit-uint32-code* [*code*]:
set-bit x n b = (*if n < 32 then uint32-set-bit x* (*integer-of-nat n*) *b else x*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint32-set-bit-code* [*code*]:
Rep-uint32 (*uint32-set-bit w n b*) =
(*if n < 0* ∨ *31 < n then Rep-uint32* (*undefined* (*set-bit* :: *uint32* ⇒ -) *w n b*)
else set-bit (*Rep-uint32 w*) (*nat-of-integer n*) *b*)
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint32-set-bit* ⇀
(*SML*) *Uint32.set'-bit* **and**
(*Haskell*) *Data'-Bits.setBitBounded* **and**
(*OCaml*) *Uint32.set'-bit* **and**
(*Scala*) *Uint32.set'-bit* **and**
(*Eval*) (*fn w => fn n => fn b => if n < 0 orelse 32 <= n then raise* (*Fail argument to uint32'-set'-bit out of bounds*) *else Uint32.set'-bit x n b*)

definition *uint32-shiftl* :: *uint32* ⇒ *integer* ⇒ *uint32*
where [*code del*]:
uint32-shiftl x n = (*if n < 0* ∨ *32* ≤ *n then undefined* (*push-bit* :: *nat* ⇒ *uint32* ⇒ -) *x n else push-bit* (*nat-of-integer n*) *x*)

lemma *shiftl-uint32-code* [*code*]: *push-bit n x* = (*if n < 32 then uint32-shiftl x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint32-shiftl-code* [*code*]:
Rep-uint32 (*uint32-shiftl w n*) =
(*if n < 0* ∨ *32* ≤ *n then Rep-uint32* (*undefined* (*push-bit* :: *nat* ⇒ *uint32* ⇒ -) *w n*) *else push-bit* (*nat-of-integer n*) (*Rep-uint32 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint32-shiftl* ⇀
(*SML*) *Uint32.shiftl* **and**
(*Haskell*) *Data'-Bits.shiftlBounded* **and**
(*OCaml*) *Uint32.shiftl* **and**
(*Scala*) *Uint32.shiftl* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 32 then raise Fail argument to uint32'-shiftl out of bounds else Uint32.shiftl x i*)

definition *uint32-shiftr* :: *uint32* ⇒ *integer* ⇒ *uint32*
where [*code del*]:
uint32-shiftr x n = (*if n < 0 ∨ 32 ≤ n then undefined* (*drop-bit* :: *nat* ⇒ *uint32* ⇒ -) *x n else drop-bit* (*nat-of-integer n*) *x*)

lemma *shiftr-uint32-code* [*code*]: *drop-bit n x* = (*if n < 32 then uint32-shiftr x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint32-shiftr-code* [*code*]:
Rep-uint32 (*uint32-shiftr w n*) =
(*if n < 0 ∨ 32 ≤ n then Rep-uint32* (*undefined* (*drop-bit* :: *nat* ⇒ *uint32* ⇒ -) *w n*) *else drop-bit* (*nat-of-integer n*) (*Rep-uint32 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint32-shiftr* ⇀
(*SML*) *Uint32.shiftr* **and**
(*Haskell*) *Data'-Bits.shiftrBounded* **and**
(*OCaml*) *Uint32.shiftr* **and**
(*Scala*) *Uint32.shiftr* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 32 then raise Fail argument to uint32'-shiftr out of bounds else Uint32.shiftr x i*)

definition *uint32-sshiftr* :: *uint32* ⇒ *integer* ⇒ *uint32*
where [*code del*]:
uint32-sshiftr x n =
(*if n < 0 ∨ 32 ≤ n then undefined signed-drop-bit-uint32 n x else signed-drop-bit-uint32* (*nat-of-integer n*) *x*)

lemma *sshiftr-uint32-code* [*code*]:
signed-drop-bit-uint32 n x =
(*if n < 32 then uint32-sshiftr x* (*integer-of-nat n*) *else if bit x 31 then* − *1 else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint32-sshiftr-code* [*code*]:
Rep-uint32 (*uint32-sshiftr w n*) =
(*if n < 0 ∨ 32 ≤ n then Rep-uint32* (*undefined signed-drop-bit-uint32 n w*) *else signed-drop-bit* (*nat-of-integer n*) (*Rep-uint32 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint32-sshiftr* ⇀
(*SML*) *Uint32.shiftr'-signed* **and**
(*Haskell*)
(*Prelude.fromInteger* (*Prelude.toInteger* (*Data'-Bits.shiftrBounded* (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint32.Int32*) -)) :: *Uint32.Word32*) **and**
(*OCaml*) *Uint32.shiftr'-signed* **and**
(*Scala*) *Uint32.shiftr'-signed* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 32 then raise Fail argument to uint32'-shiftr'-signed out of bounds else Uint32.shiftr'-signed x i*)

context
includes *bit-operations-syntax*
begin

lemma *uint32-msb-test-bit*: *msb x* ⟷ *bit* (*x* :: *uint32*) *31*
⟨*proof*⟩

lemma *msb-uint32-code* [*code*]: *msb x* ⟷ *uint32-test-bit x 31*
⟨*proof*⟩

lemma *uint32-of-int-code* [*code*]:
uint32-of-int i = *Uint32* (*integer-of-int i*)
including *integer.lifting* ⟨*proof*⟩

lemma *int-of-uint32-code* [*code*]:
int-of-uint32 x = *int-of-integer* (*integer-of-uint32 x*)
including *integer.lifting* ⟨*proof*⟩

lemma *uint32-of-nat-code* [*code*]:
uint32-of-nat = *uint32-of-int* ∘ *int*
⟨*proof*⟩

lemma *nat-of-uint32-code* [*code*]:
nat-of-uint32 x = *nat-of-integer* (*integer-of-uint32 x*)
⟨*proof*⟩ **including** *integer.lifting* ⟨*proof*⟩

definition *integer-of-uint32-signed* :: *uint32* ⇒ *integer*
where
integer-of-uint32-signed n = (*if bit n 31 then undefined integer-of-uint32 n else integer-of-uint32 n*)

lemma *integer-of-uint32-signed-code* [*code*]:
integer-of-uint32-signed n =
(*if bit n 31 then undefined integer-of-uint32 n else integer-of-int* (*uint* (*Rep-uint32' n*)))
⟨*proof*⟩

lemma *integer-of-uint32-code* [*code*]:
integer-of-uint32 n =

(*if bit n 31 then integer-of-uint32-signed* (*n AND 0x7FFFFFFF*) *OR 0x80000000 else integer-of-uint32-signed n*)
⟨*proof*⟩
including *integer.lifting* ⟨*proof*⟩

end

code-printing
constant *integer-of-uint32* ⇀
(*SML*) *IntInf.fromLarge* (*Word32.toLargeInt* -) : *IntInf.int* **and**
(*Haskell*) *Prelude.toInteger*
| **constant** *integer-of-uint32-signed* ⇀
(*OCaml*) *Z.of'-int32* **and**
(*Scala*) *BigInt*

7.3 Quickcheck setup

definition *uint32-of-natural* :: *natural* ⇒ *uint32*
where *uint32-of-natural x* ≡ *Uint32* (*integer-of-natural x*)

instantiation *uint32* :: {*random*, *exhaustive*, *full-exhaustive*} **begin**
definition *random-uint32* ≡ *qc-random-cnv uint32-of-natural*
definition *exhaustive-uint32* ≡ *qc-exhaustive-cnv uint32-of-natural*
definition *full-exhaustive-uint32* ≡ *qc-full-exhaustive-cnv uint32-of-natural*
instance ⟨*proof*⟩
end

instantiation *uint32* :: *narrowing* **begin**

interpretation *quickcheck-narrowing-samples*
λi. let x = Uint32 i in (*x*, *0xFFFFFFFF − x*) *0*
Typerep.Typerep (*STR ''Uint32.uint32''*) [] ⟨*proof*⟩

definition *narrowing-uint32 d = qc-narrowing-drawn-from* (*narrowing-samples d*) *d*
declare [[*code drop*: *partial-term-of* :: *uint32 itself* ⇒ -]]
lemmas *partial-term-of-uint32* [*code*] = *partial-term-of-code*

instance ⟨*proof*⟩
end

end

Chapter 8

Unsigned words of 16 bits

theory *Uint16* **imports**
Code-Target-Word-Base Word-Type-Copies
begin

Restriction for ML code generation: This theory assumes that the ML system provides a Word16 implementation (mlton does, but PolyML 5.5 does not). Therefore, the code setup lives in the target *SML-word* rather than *SML*. This ensures that code generation still works as long as *uint16* is not involved. For the target *SML* itself, no special code generation for this type is set up. Nevertheless, it should work by emulation via *16 word* if the theory *Code-Target-Bits-Int* is imported.

Restriction for OCaml code generation: OCaml does not provide an int16 type, so no special code generation for this type is set up.

8.1 Type definition and primitive operations

typedef *uint16* = ‹*UNIV* :: *16 word set*› ⟨*proof*⟩

global-interpretation *uint16*: *word-type-copy Abs-uint16 Rep-uint16*
⟨*proof*⟩

setup-lifting *type-definition-uint16*

declare *uint16.of-word-of* [*code abstype*]

declare *Quotient-uint16* [*transfer-rule*]

instantiation *uint16* :: ‹{*comm-ring-1*, *semiring-modulo*, *equal*, *linorder*}›
begin

lift-definition *zero-uint16* :: *uint16* **is** *0* ⟨*proof*⟩
lift-definition *one-uint16* :: *uint16* **is** *1* ⟨*proof*⟩
lift-definition *plus-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹(+)› ⟨*proof*⟩

lift-definition *uminus-uint16* :: ‹*uint16* ⇒ *uint16*› **is** *uminus* ⟨*proof*⟩
lift-definition *minus-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹(−)› ⟨*proof*⟩
lift-definition *times-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹(∗)› ⟨*proof*⟩
lift-definition *divide-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹(*div*)› ⟨*proof*⟩
lift-definition *modulo-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹(*mod*)› ⟨*proof*⟩
lift-definition *equal-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *bool*› **is** ‹*HOL.equal*› ⟨*proof*⟩
lift-definition *less-eq-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *bool*› **is** ‹(≤)› ⟨*proof*⟩
lift-definition *less-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *bool*› **is** ‹(<)› ⟨*proof*⟩

global-interpretation *uint16*: *word-type-copy-ring Abs-uint16 Rep-uint16*
⟨*proof*⟩

instance ⟨*proof*⟩

end

instantiation *uint16* :: *ring-bit-operations*
begin

lift-definition *bit-uint16* :: ‹*uint16* ⇒ *nat* ⇒ *bool*› **is** *bit* ⟨*proof*⟩
lift-definition *not-uint16* :: ‹*uint16* ⇒ *uint16*› **is** ‹*Bit-Operations.not*› ⟨*proof*⟩
lift-definition *and-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹*Bit-Operations.and*›
⟨*proof*⟩
lift-definition *or-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹*Bit-Operations.or*›
⟨*proof*⟩
lift-definition *xor-uint16* :: ‹*uint16* ⇒ *uint16* ⇒ *uint16*› **is** ‹*Bit-Operations.xor*›
⟨*proof*⟩
lift-definition *mask-uint16* :: ‹*nat* ⇒ *uint16*› **is** *mask* ⟨*proof*⟩
lift-definition *push-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *push-bit* ⟨*proof*⟩
lift-definition *drop-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *drop-bit* ⟨*proof*⟩
lift-definition *signed-drop-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *signed-drop-bit*
⟨*proof*⟩
lift-definition *take-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *take-bit* ⟨*proof*⟩
lift-definition *set-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *Bit-Operations.set-bit*
⟨*proof*⟩
lift-definition *unset-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *unset-bit* ⟨*proof*⟩
lift-definition *flip-bit-uint16* :: ‹*nat* ⇒ *uint16* ⇒ *uint16*› **is** *flip-bit* ⟨*proof*⟩

global-interpretation *uint16*: *word-type-copy-bits Abs-uint16 Rep-uint16 signed-drop-bit-uint16*
⟨*proof*⟩

instance
⟨*proof*⟩

end

lift-definition *uint16-of-nat* :: ‹*nat* ⇒ *uint16*›
is *word-of-nat* ⟨*proof*⟩

lift-definition *nat-of-uint16* :: ‹*uint16* ⇒ *nat*›
  is *unat* ⟨*proof*⟩

lift-definition *uint16-of-int* :: ‹*int* ⇒ *uint16*›
  is *word-of-int* ⟨*proof*⟩

lift-definition *int-of-uint16* :: ‹*uint16* ⇒ *int*›
  is *uint* ⟨*proof*⟩

context
  includes *integer.lifting*
begin

lift-definition *Uint16* :: ‹*integer* ⇒ *uint16*›
  is *word-of-int* ⟨*proof*⟩

lift-definition *integer-of-uint16* :: ‹*uint16* ⇒ *integer*›
  is *uint* ⟨*proof*⟩

end

global-interpretation *uint16*: *word-type-copy-more Abs-uint16 Rep-uint16 signed-drop-bit-uint16 uint16-of-nat nat-of-uint16 uint16-of-int int-of-uint16 Uint16 integer-of-uint16*
  ⟨*proof*⟩

instantiation *uint16* :: {*size*, *msb*, *lsb*, *set-bit*, *bit-comprehension*}
begin

lift-definition *size-uint16* :: ‹*uint16* ⇒ *nat*› **is** *size* ⟨*proof*⟩

lift-definition *msb-uint16* :: ‹*uint16* ⇒ *bool*› **is** *msb* ⟨*proof*⟩
lift-definition *lsb-uint16* :: ‹*uint16* ⇒ *bool*› **is** *lsb* ⟨*proof*⟩

Workaround: avoid name space clash by spelling out *lift-definition* explicitly.

definition *set-bit-uint16* :: ‹*uint16* ⇒ *nat* ⇒ *bool* ⇒ *uint16*›
  where *set-bit-uint16-eq*: ‹*set-bit-uint16 a n b* = (*if b then Bit-Operations.set-bit else unset-bit*) *n a*›

context
  includes *lifting-syntax*
begin

lemma *set-bit-uint16-transfer* [*transfer-rule*]:
  ‹(*cr-uint16* ===> (=) ===> (⟷) ===> *cr-uint16*) *Generic-set-bit.set-bit Generic-set-bit.set-bit*›
  ⟨*proof*⟩

end

lift-definition *set-bits-uint16* :: ‹(*nat* ⇒ *bool*) ⇒ *uint16*› **is** *set-bits* ⟨*proof*⟩
lift-definition *set-bits-aux-uint16* :: ‹(*nat* ⇒ *bool*) ⇒ *nat* ⇒ *uint16* ⇒ *uint16*› **is**
set-bits-aux ⟨*proof*⟩

global-interpretation *uint16*: *word-type-copy-misc Abs-uint16 Rep-uint16 signed-drop-bit-uint16*
uint16-of-nat nat-of-uint16 uint16-of-int int-of-uint16 Uint16 integer-of-uint16 16
set-bits-aux-uint16
⟨*proof*⟩

instance ⟨*proof*⟩

end

8.2 Code setup

code-printing code-module *Uint16* ⇀ (*SML-word*)
‹(∗ *Test that words can handle numbers between 0 and 15* ∗)
val - = if 4 <= Word.wordSize then () else raise (Fail (wordSize less than 4));

structure Uint16 : sig
val set-bit : Word16.word -> IntInf.int -> bool -> Word16.word
val shiftl : Word16.word -> IntInf.int -> Word16.word
val shiftr : Word16.word -> IntInf.int -> Word16.word
val shiftr-signed : Word16.word -> IntInf.int -> Word16.word
val test-bit : Word16.word -> IntInf.int -> bool
end = struct

fun set-bit x n b =
let val mask = Word16.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))
in if b then Word16.orb (x, mask)
else Word16.andb (x, Word16.notb mask)
end

fun shiftl x n =
Word16.<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
Word16.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
Word16.~>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
Word16.andb (x, Word16.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n)))
<> Word16.fromInt 0

end; (∗ *struct Uint16* ∗)›
code-reserved *SML-word Uint16*

code-printing code-module *Uint16* ⇀ (*Haskell*)
‹*module Uint16*(*Int16*, *Word16*) *where*

import Data.Int(*Int16*)
import Data.Word(*Word16*)›
code-reserved *Haskell Uint16*

Scala provides unsigned 16-bit numbers as Char.

code-printing code-module *Uint16* ⇀ (*Scala*)
‹*object Uint16* {

def set-bit(*x*: *scala.Char*, *n*: *BigInt*, *b*: *Boolean*) : *scala.Char* =
if (*b*)
(*x* | (*1.toChar* << *n.intValue*)).*toChar*
else
(*x* & (*1.toChar* << *n.intValue*).*unary-*~).*toChar*

def shiftl(*x*: *scala.Char*, *n*: *BigInt*) : *scala.Char* = (*x* << *n.intValue*).*toChar*

def shiftr(*x*: *scala.Char*, *n*: *BigInt*) : *scala.Char* = (*x* >>> *n.intValue*).*toChar*

def shiftr-signed(*x*: *scala.Char*, *n*: *BigInt*) : *scala.Char* = (*x.toShort* >> *n.intValue*).*toChar*

def test-bit(*x*: *scala.Char*, *n*: *BigInt*) : *Boolean* = (*x* & (*1.toChar* << *n.intValue*))
!= *0*

} /∗ *object Uint16* ∗/›
code-reserved *Scala Uint16*

Avoid *Abs-uint16* in generated code, use *Rep-uint16′* instead. The symbolic implementations for code_simp use *Rep-uint16*.

The new destructor *Rep-uint16′* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint16* directly, because that makes code_simp loop.

If code generation raises Match, some equation probably contains *Rep-uint16* ([code abstract] equations for *uint16* may use *Rep-uint16* because these instances will be folded away.)

To convert *16 word* values into *uint16*, use *Abs-uint16′*.

definition *Rep-uint16′* **where** [*simp*]: *Rep-uint16′* = *Rep-uint16*

lemma *Rep-uint16′-transfer* [*transfer-rule*]:
rel-fun cr-uint16 (=) (λ*x*. *x*) *Rep-uint16′*
⟨*proof*⟩

lemma *Rep-uint16′-code* [*code*]: *Rep-uint16′ x* = (*BITS n. bit x n*)
⟨*proof*⟩

lift-definition *Abs-uint16'* :: *16 word* ⇒ *uint16* **is** λ*x* :: *16 word*. *x* ⟨*proof*⟩

lemma *Abs-uint16'-code* [*code*]:
Abs-uint16' x = *Uint16* (*integer-of-int* (*uint x*))
including *integer.lifting* ⟨*proof*⟩

declare [[*code drop*: *term-of-class.term-of* :: *uint16* ⇒ -]]

lemma *term-of-uint16-code* [*code*]:
defines *TR* ≡ *typerep.Typerep* **and** *bit0* ≡ *STR ''Numeral-Type.bit0''* **shows**
term-of-class.term-of x =
Code-Evaluation.App (*Code-Evaluation.Const* (*STR ''Uint16.uint16.Abs-uint16''*)
(*TR* (*STR ''fun''*) [*TR* (*STR ''Word.word''*) [*TR bit0* [*TR bit0* [*TR bit0* [*TR bit0*
[*TR* (*STR ''Numeral-Type.num1''*) []]]]], *TR* (*STR ''Uint16.uint16''*) []])) (*term-of-class.term-of* (*Rep-uint16' x*))
⟨*proof*⟩

lemma *Uint16-code* [*code*]: *Rep-uint16* (*Uint16 i*) = *word-of-int* (*int-of-integer-symbolic i*)
⟨*proof*⟩

code-printing
type-constructor *uint16* ⇀
(*SML-word*) *Word16.word* **and**
(*Haskell*) *Uint16.Word16* **and**
(*Scala*) *scala.Char*
| **constant** *Uint16* ⇀
(*SML-word*) *Word16.fromLargeInt* (*IntInf.toLarge* -) **and**
(*Haskell*) (*Prelude.fromInteger* - :: *Uint16.Word16*) **and**
(*Haskell-Quickcheck*) (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint16.Word16*)
and
(*Scala*) -.*charValue*
| **constant** *0* :: *uint16* ⇀
(*SML-word*) (*Word16.fromInt 0*) **and**
(*Haskell*) (*0* :: *Uint16.Word16*) **and**
(*Scala*) *0*
| **constant** *1* :: *uint16* ⇀
(*SML-word*) (*Word16.fromInt 1*) **and**
(*Haskell*) (*1* :: *Uint16.Word16*) **and**
(*Scala*) *1*
| **constant** *plus* :: *uint16* ⇒ - ⇒ - ⇀
(*SML-word*) *Word16.+* ((-), (-)) **and**
(*Haskell*) **infixl** *6* + **and**
(*Scala*) (- +/ -).*toChar*
| **constant** *uminus* :: *uint16* ⇒ - ⇀
(*SML-word*) *Word16.~* **and**
(*Haskell*) *negate* **and**
(*Scala*) (− -).*toChar*
| **constant** *minus* :: *uint16* ⇒ - ⇀

(*SML-word*) *Word16.−* ((-), (-)) **and**
(*Haskell*) **infixl** *6 −* **and**
(*Scala*) (- −/ -).*toChar*
| **constant** *times* :: *uint16* ⇒ - ⇒ - ⇀
(*SML-word*) *Word16.∗* ((-), (-)) **and**
(*Haskell*) **infixl** *7 ∗* **and**
(*Scala*) (- ∗/ -).*toChar*
| **constant** *HOL.equal* :: *uint16* ⇒ - ⇒ *bool* ⇀
(*SML-word*) !((- : *Word16.word*) = -) **and**
(*Haskell*) **infix** *4 ==* **and**
(*Scala*) **infixl** *5 ==*
| **class-instance** *uint16* :: *equal* ⇀ (*Haskell*) −
| **constant** *less-eq* :: *uint16* ⇒ - ⇒ *bool* ⇀
(*SML-word*) *Word16.<=* ((-), (-)) **and**
(*Haskell*) **infix** *4 <=* **and**
(*Scala*) **infixl** *4 <=*
| **constant** *less* :: *uint16* ⇒ - ⇒ *bool* ⇀
(*SML-word*) *Word16.<* ((-), (-)) **and**
(*Haskell*) **infix** *4 <* **and**
(*Scala*) **infixl** *4 <*
| **constant** *Bit-Operations.not* :: *uint16* ⇒ - ⇀
(*SML-word*) *Word16.notb* **and**
(*Haskell*) *Data′-Bits.complement* **and**
(*Scala*) -.*unary′-*~.*toChar*
| **constant** *Bit-Operations.and* :: *uint16* ⇒ - ⇀
(*SML-word*) *Word16.andb* ((-),/ (-)) **and**
(*Haskell*) **infixl** *7 Data-Bits..&.* **and**
(*Scala*) (- & -).*toChar*
| **constant** *Bit-Operations.or* :: *uint16* ⇒ - ⇀
(*SML-word*) *Word16.orb* ((-),/ (-)) **and**
(*Haskell*) **infixl** *5 Data-Bits..|.* **and**
(*Scala*) (- | -).*toChar*
| **constant** *Bit-Operations.xor* :: *uint16* ⇒ - ⇀
(*SML-word*) *Word16.xorb* ((-),/ (-)) **and**
(*Haskell*) *Data′-Bits.xor* **and**
(*Scala*) (- ^ -).*toChar*

definition *uint16-div* :: *uint16* ⇒ *uint16* ⇒ *uint16*
where *uint16-div x y* = (*if y* = *0 then undefined* ((*div*) :: *uint16* ⇒ -) *x* (*0* :: *uint16*) *else x div y*)

definition *uint16-mod* :: *uint16* ⇒ *uint16* ⇒ *uint16*
where *uint16-mod x y* = (*if y* = *0 then undefined* ((*mod*) :: *uint16* ⇒ -) *x* (*0* :: *uint16*) *else x mod y*)

context includes *undefined-transfer* **begin**

lemma *div-uint16-code* [*code*]: *x div y* = (*if y* = *0 then 0 else uint16-div x y*)
⟨*proof*⟩

lemma *mod-uint16-code* [*code*]: *x mod y* = (*if y* = *0 then x else uint16-mod x y*)
⟨*proof*⟩

lemma *uint16-div-code* [*code*]:
Rep-uint16 (*uint16-div x y*) =
(*if y* = *0 then Rep-uint16* (*undefined* ((*div*) :: *uint16* ⇒ -) *x* (*0* :: *uint16*)) *else Rep-uint16 x div Rep-uint16 y*)
⟨*proof*⟩

lemma *uint16-mod-code* [*code*]:
Rep-uint16 (*uint16-mod x y*) =
(*if y* = *0 then Rep-uint16* (*undefined* ((*mod*) :: *uint16* ⇒ -) *x* (*0* :: *uint16*)) *else Rep-uint16 x mod Rep-uint16 y*)
⟨*proof*⟩

end

code-printing constant *uint16-div* ⇀
(*SML-word*) *Word16.div* ((-), (-)) **and**
(*Haskell*) *Prelude.div* **and**
(*Scala*) (- '/ -).*toChar*
| **constant** *uint16-mod* ⇀
(*SML-word*) *Word16.mod* ((-), (-)) **and**
(*Haskell*) *Prelude.mod* **and**
(*Scala*) (- % -).*toChar*

definition *uint16-test-bit* :: *uint16* ⇒ *integer* ⇒ *bool*
where [*code del*]:
uint16-test-bit x n =
(*if n* < *0* ∨ *15* < *n then undefined* (*bit* :: *uint16* ⇒ -) *x n*
else bit x (*nat-of-integer n*))

lemma *test-bit-uint16-code* [*code*]:
bit x n ⟷ *n* < *16* ∧ *uint16-test-bit x* (*integer-of-nat n*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint16-test-bit-code* [*code*]:
uint16-test-bit w n =
(*if n* < *0* ∨ *15* < *n then undefined* (*bit* :: *uint16* ⇒ -) *w n else bit* (*Rep-uint16 w*) (*nat-of-integer n*))
⟨*proof*⟩

code-printing constant *uint16-test-bit* ⇀
(*SML-word*) *Uint16.test'-bit* **and**
(*Haskell*) *Data'-Bits.testBitBounded* **and**
(*Scala*) *Uint16.test'-bit*

definition *uint16-set-bit* :: *uint16* ⇒ *integer* ⇒ *bool* ⇒ *uint16*

where [*code del*]:
uint16-set-bit x n b =
(*if n* < *0* ∨ *15* < *n then undefined* (*set-bit* :: *uint16* ⇒ -) *x n b*
else set-bit x (*nat-of-integer n*) *b*)

lemma *set-bit-uint16-code* [*code*]:
set-bit x n b = (*if n* < *16 then uint16-set-bit x* (*integer-of-nat n*) *b else x*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint16-set-bit-code* [*code*]:
Rep-uint16 (*uint16-set-bit w n b*) =
(*if n* < *0* ∨ *15* < *n then Rep-uint16* (*undefined* (*set-bit* :: *uint16* ⇒ -) *w n b*)
else set-bit (*Rep-uint16 w*) (*nat-of-integer n*) *b*)
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint16-set-bit* ⇀
(*SML-word*) *Uint16.set'-bit* **and**
(*Haskell*) *Data'-Bits.setBitBounded* **and**
(*Scala*) *Uint16.set'-bit*

definition *uint16-shiftl* :: *uint16* ⇒ *integer* ⇒ *uint16*
where [*code del*]:
uint16-shiftl x n = (*if n* < *0* ∨ *16* ≤ *n then undefined* (*push-bit* :: *nat* ⇒ *uint16* ⇒ -) *x n else push-bit* (*nat-of-integer n*) *x*)

lemma *shiftl-uint16-code* [*code*]: *push-bit n x* = (*if n* < *16 then uint16-shiftl x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint16-shiftl-code* [*code*]:
Rep-uint16 (*uint16-shiftl w n*) =
(*if n* < *0* ∨ *16* ≤ *n then Rep-uint16* (*undefined* (*push-bit* :: *nat* ⇒ *uint16* ⇒ -) *w n*)
else push-bit (*nat-of-integer n*) (*Rep-uint16 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint16-shiftl* ⇀
(*SML-word*) *Uint16.shiftl* **and**
(*Haskell*) *Data'-Bits.shiftlBounded* **and**
(*Scala*) *Uint16.shiftl*

definition *uint16-shiftr* :: *uint16* ⇒ *integer* ⇒ *uint16*
where [*code del*]:
uint16-shiftr x n = (*if n* < *0* ∨ *16* ≤ *n then undefined* (*drop-bit* :: *nat* ⇒ *uint16* ⇒ -) *x n else drop-bit* (*nat-of-integer n*) *x*)

lemma *shiftr-uint16-code* [*code*]: *drop-bit n x* = (*if n* < *16 then uint16-shiftr x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint16-shiftr-code* [*code*]:
Rep-uint16 (*uint16-shiftr w n*) =
(*if n* < *0* ∨ *16* ≤ *n then Rep-uint16* (*undefined* (*drop-bit* :: *nat* ⇒ *uint16* ⇒ -)
w n)
else drop-bit (*nat-of-integer n*) (*Rep-uint16 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint16-shiftr* ⇀
(*SML-word*) *Uint16.shiftr* **and**
(*Haskell*) *Data′-Bits.shiftrBounded* **and**
(*Scala*) *Uint16.shiftr*

definition *uint16-sshiftr* :: *uint16* ⇒ *integer* ⇒ *uint16*
where [*code del*]:
uint16-sshiftr x n =
(*if n* < *0* ∨ *16* ≤ *n then undefined signed-drop-bit-uint16 n x else signed-drop-bit-uint16*
(*nat-of-integer n*) *x*)

lemma *sshiftr-uint16-code* [*code*]:
signed-drop-bit-uint16 n x =
(*if n* < *16 then uint16-sshiftr x* (*integer-of-nat n*) *else if bit x 15 then* −*1 else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint16-sshiftr-code* [*code*]:
Rep-uint16 (*uint16-sshiftr w n*) =
(*if n* < *0* ∨ *16* ≤ *n then Rep-uint16* (*undefined signed-drop-bit-uint16 n w*)
else signed-drop-bit (*nat-of-integer n*) (*Rep-uint16 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint16-sshiftr* ⇀
(*SML-word*) *Uint16.shiftr′-signed* **and**
(*Haskell*)
(*Prelude.fromInteger* (*Prelude.toInteger* (*Data′-Bits.shiftrBounded* (*Prelude.fromInteger*
(*Prelude.toInteger* -) :: *Uint16.Int16*) -)) :: *Uint16.Word16*) **and**
(*Scala*) *Uint16.shiftr′-signed*

lemma *uint16-msb-test-bit*: *msb x* ⟷ *bit* (*x* :: *uint16*) *15*
⟨*proof*⟩

lemma *msb-uint16-code* [*code*]: *msb x* ⟷ *uint16-test-bit x 15*
⟨*proof*⟩

lemma *uint16-of-int-code* [*code*]: *uint16-of-int i* = *Uint16* (*integer-of-int i*)
including *integer.lifting* ⟨*proof*⟩

lemma *int-of-uint16-code* [*code*]:
int-of-uint16 x = *int-of-integer* (*integer-of-uint16 x*)
⟨*proof*⟩

lemma *uint16-of-nat-code* [*code*]:
uint16-of-nat = *uint16-of-int* ∘ *int*
⟨*proof*⟩

lemma *nat-of-uint16-code* [*code*]:
nat-of-uint16 x = *nat-of-integer* (*integer-of-uint16 x*)
⟨*proof*⟩ **including** *integer.lifting* ⟨*proof*⟩

lemma *integer-of-uint16-code* [*code*]:
integer-of-uint16 n = *integer-of-int* (*uint* (*Rep-uint16′ n*))
⟨*proof*⟩

code-printing
constant *integer-of-uint16* ⇀
(*SML-word*) *Word16.toInt* - : *IntInf.int* **and**
(*Haskell*) *Prelude.toInteger* **and**
(*Scala*) *BigInt*

8.3 Quickcheck setup

definition *uint16-of-natural* :: *natural* ⇒ *uint16*
where *uint16-of-natural x* ≡ *Uint16* (*integer-of-natural x*)

instantiation *uint16* :: {*random*, *exhaustive*, *full-exhaustive*} **begin**
definition *random-uint16* ≡ *qc-random-cnv uint16-of-natural*
definition *exhaustive-uint16* ≡ *qc-exhaustive-cnv uint16-of-natural*
definition *full-exhaustive-uint16* ≡ *qc-full-exhaustive-cnv uint16-of-natural*
instance ⟨*proof*⟩
end

instantiation *uint16* :: *narrowing* **begin**

interpretation *quickcheck-narrowing-samples*
λi. let x = *Uint16 i in* (*x*, *0xFFFF* − *x*) *0*
Typerep.Typerep (*STR ″Uint16.uint16″*) [] ⟨*proof*⟩

definition *narrowing-uint16 d* = *qc-narrowing-drawn-from* (*narrowing-samples d*)
d
declare [[*code drop*: *partial-term-of* :: *uint16 itself* ⇒ -]]
lemmas *partial-term-of-uint16* [*code*] = *partial-term-of-code*

instance ⟨*proof*⟩
end

end

Chapter 9

Unsigned words of 8 bits

theory *Uint8* **imports**
Code-Target-Word-Base Word-Type-Copies
begin

Restriction for OCaml code generation: OCaml does not provide an int8 type, so no special code generation for this type is set up. If the theory *Code-Target-Bits-Int* is imported, the type *uint8* is emulated via *8 word*.

9.1 Type definition and primitive operations

typedef *uint8* = ‹*UNIV* :: *8 word set*› ⟨*proof*⟩

global-interpretation *uint8*: *word-type-copy Abs-uint8 Rep-uint8*
⟨*proof*⟩

setup-lifting *type-definition-uint8*

declare *uint8.of-word-of* [*code abstype*]

declare *Quotient-uint8* [*transfer-rule*]

instantiation *uint8* :: ‹{*comm-ring-1*, *semiring-modulo*, *equal*, *linorder*}›
begin

lift-definition *zero-uint8* :: *uint8* **is** *0* ⟨*proof*⟩
lift-definition *one-uint8* :: *uint8* **is** *1* ⟨*proof*⟩
lift-definition *plus-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹(+)› ⟨*proof*⟩
lift-definition *uminus-uint8* :: ‹*uint8* ⇒ *uint8*› **is** *uminus* ⟨*proof*⟩
lift-definition *minus-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹(−)› ⟨*proof*⟩
lift-definition *times-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹(∗)› ⟨*proof*⟩
lift-definition *divide-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹(*div*)› ⟨*proof*⟩
lift-definition *modulo-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹(*mod*)› ⟨*proof*⟩
lift-definition *equal-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *bool*› **is** ‹*HOL.equal*› ⟨*proof*⟩
lift-definition *less-eq-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *bool*› **is** ‹(≤)› ⟨*proof*⟩

lift-definition *less-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *bool*› **is** ‹(<)› ⟨*proof*⟩

global-interpretation *uint8*: *word-type-copy-ring Abs-uint8 Rep-uint8*
⟨*proof*⟩

instance ⟨*proof*⟩

end

instantiation *uint8* :: *ring-bit-operations*
begin

lift-definition *bit-uint8* :: ‹*uint8* ⇒ *nat* ⇒ *bool*› **is** *bit* ⟨*proof*⟩
lift-definition *not-uint8* :: ‹*uint8* ⇒ *uint8*› **is** ‹*Bit-Operations.not*› ⟨*proof*⟩
lift-definition *and-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹*Bit-Operations.and*›
⟨*proof*⟩
lift-definition *or-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹*Bit-Operations.or*› ⟨*proof*⟩
lift-definition *xor-uint8* :: ‹*uint8* ⇒ *uint8* ⇒ *uint8*› **is** ‹*Bit-Operations.xor*› ⟨*proof*⟩
lift-definition *mask-uint8* :: ‹*nat* ⇒ *uint8*› **is** *mask* ⟨*proof*⟩
lift-definition *push-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *push-bit* ⟨*proof*⟩
lift-definition *drop-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *drop-bit* ⟨*proof*⟩
lift-definition *signed-drop-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *signed-drop-bit*
⟨*proof*⟩
lift-definition *take-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *take-bit* ⟨*proof*⟩
lift-definition *set-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *Bit-Operations.set-bit*
⟨*proof*⟩
lift-definition *unset-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *unset-bit* ⟨*proof*⟩
lift-definition *flip-bit-uint8* :: ‹*nat* ⇒ *uint8* ⇒ *uint8*› **is** *flip-bit* ⟨*proof*⟩

global-interpretation *uint8*: *word-type-copy-bits Abs-uint8 Rep-uint8 signed-drop-bit-uint8*
⟨*proof*⟩

instance
⟨*proof*⟩

end

lift-definition *uint8-of-nat* :: ‹*nat* ⇒ *uint8*›
is *word-of-nat* ⟨*proof*⟩

lift-definition *nat-of-uint8* :: ‹*uint8* ⇒ *nat*›
is *unat* ⟨*proof*⟩

lift-definition *uint8-of-int* :: ‹*int* ⇒ *uint8*›
is *word-of-int* ⟨*proof*⟩

lift-definition *int-of-uint8* :: ‹*uint8* ⇒ *int*›
is *uint* ⟨*proof*⟩

context
includes *integer.lifting*
begin

lift-definition *Uint8* :: ‹*integer* ⇒ *uint8*›
is *word-of-int* ⟨*proof*⟩

lift-definition *integer-of-uint8* :: ‹*uint8* ⇒ *integer*›
is *uint* ⟨*proof*⟩

end

global-interpretation *uint8*: *word-type-copy-more Abs-uint8 Rep-uint8 signed-drop-bit-uint8 uint8-of-nat nat-of-uint8 uint8-of-int int-of-uint8 Uint8 integer-of-uint8*
⟨*proof*⟩

instantiation *uint8* :: {*size*, *msb*, *lsb*, *set-bit*, *bit-comprehension*}
begin

lift-definition *size-uint8* :: ‹*uint8* ⇒ *nat*› **is** *size* ⟨*proof*⟩

lift-definition *msb-uint8* :: ‹*uint8* ⇒ *bool*› **is** *msb* ⟨*proof*⟩
lift-definition *lsb-uint8* :: ‹*uint8* ⇒ *bool*› **is** *lsb* ⟨*proof*⟩

Workaround: avoid name space clash by spelling out *lift-definition* explicitly.

definition *set-bit-uint8* :: ‹*uint8* ⇒ *nat* ⇒ *bool* ⇒ *uint8*›
where *set-bit-uint8-eq*: ‹*set-bit-uint8 a n b* = (*if b then Bit-Operations.set-bit else unset-bit*) *n a*›

context
includes *lifting-syntax*
begin

lemma *set-bit-uint8-transfer* [*transfer-rule*]:
‹(*cr-uint8* ===> (=) ===> (⟷) ===> *cr-uint8*) *Generic-set-bit.set-bit Generic-set-bit.set-bit*›
⟨*proof*⟩

end

lift-definition *set-bits-uint8* :: ‹(*nat* ⇒ *bool*) ⇒ *uint8*› **is** *set-bits* ⟨*proof*⟩
lift-definition *set-bits-aux-uint8* :: ‹(*nat* ⇒ *bool*) ⇒ *nat* ⇒ *uint8* ⇒ *uint8*› **is** *set-bits-aux* ⟨*proof*⟩

global-interpretation *uint8*: *word-type-copy-misc Abs-uint8 Rep-uint8 signed-drop-bit-uint8 uint8-of-nat nat-of-uint8 uint8-of-int int-of-uint8 Uint8 integer-of-uint8 8 set-bits-aux-uint8*
⟨*proof*⟩

instance ⟨*proof*⟩

end

9.2 Code setup

code-printing code-module *Uint8* ⇀ (*SML*)
‹(∗ *Test that words can handle numbers between 0 and 3* ∗)
val - = if 3 <= Word.wordSize then () else raise (Fail (wordSize less than 3));

structure Uint8 : sig
val set-bit : Word8.word -> IntInf.int -> bool -> Word8.word
val shiftl : Word8.word -> IntInf.int -> Word8.word
val shiftr : Word8.word -> IntInf.int -> Word8.word
val shiftr-signed : Word8.word -> IntInf.int -> Word8.word
val test-bit : Word8.word -> IntInf.int -> bool
end = struct

fun set-bit x n b =
let val mask = Word8.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))
in if b then Word8.orb (x, mask)
else Word8.andb (x, Word8.notb mask)
end

fun shiftl x n =
Word8.<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
Word8.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
Word8.~>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
Word8.andb (x, Word8.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))) <>
Word8.fromInt 0

end; (∗ *struct Uint8* ∗)›
code-reserved *SML Uint8*

code-printing code-module *Uint8* ⇀ (*Haskell*)
‹*module Uint8*(*Int8*, *Word8*) *where*

import Data.Int(*Int8*)
import Data.Word(*Word8*)›
code-reserved *Haskell Uint8*

Scala provides only signed 8bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

code-printing code-module *Uint8* ⇀ (*Scala*)

‹*object Uint8* {

def less(*x*: *Byte*, *y*: *Byte*) : *Boolean* =
  if (*x* < *0*) *y* < *0* &&& *x* < *y*
  else y < *0* || *x* < *y*

def less-eq(*x*: *Byte*, *y*: *Byte*) : *Boolean* =
  if (*x* < *0*) *y* < *0* &&& *x* <= *y*
  else y < *0* || *x* <= *y*

def set-bit(*x*: *Byte*, *n*: *BigInt*, *b*: *Boolean*) : *Byte* =
  if (*b*)
    (*x* | (*1* << *n.intValue*)).*toByte*
  else
    (*x* & (*1* << *n.intValue*).*unary-*~).*toByte*

def shiftl(*x*: *Byte*, *n*: *BigInt*) : *Byte* = (*x* << *n.intValue*).*toByte*

def shiftr(*x*: *Byte*, *n*: *BigInt*) : *Byte* = ((*x* & *255*) >>> *n.intValue*).*toByte*

def shiftr-signed(*x*: *Byte*, *n*: *BigInt*) : *Byte* = (*x* >> *n.intValue*).*toByte*

def test-bit(*x*: *Byte*, *n*: *BigInt*) : *Boolean* =
  (*x* & (*1* << *n.intValue*)) != *0*

} /* *object Uint8* */›
code-reserved *Scala Uint8*

Avoid *Abs-uint8* in generated code, use *Rep-uint8'* instead. The symbolic implementations for code_simp use *Rep-uint8*.

The new destructor *Rep-uint8'* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint8* directly, because that makes code_simp loop.

If code generation raises Match, some equation probably contains *Rep-uint8* ([code abstract] equations for *uint8* may use *Rep-uint8* because these instances will be folded away.)

To convert *8 word* values into *uint8*, use *Abs-uint8'*.

definition *Rep-uint8'* **where** [*simp*]: *Rep-uint8'* = *Rep-uint8*

lemma *Rep-uint8'-transfer* [*transfer-rule*]:
  rel-fun cr-uint8 (=) (λ*x*. *x*) *Rep-uint8'*
⟨*proof*⟩

lemma *Rep-uint8'-code* [*code*]: *Rep-uint8' x* = (*BITS n. bit x n*)
  ⟨*proof*⟩

lift-definition *Abs-uint8'* :: *8 word* ⇒ *uint8* **is** λ*x* :: *8 word*. *x* ⟨*proof*⟩

lemma *Abs-uint8'-code* [*code*]: *Abs-uint8' x* = *Uint8* (*integer-of-int* (*uint x*))
including *integer.lifting* ⟨*proof*⟩

declare [[*code drop*: *term-of-class.term-of* :: *uint8* ⇒ -]]

lemma *term-of-uint8-code* [*code*]:
defines *TR* ≡ *typerep.Typerep* **and** *bit0* ≡ *STR ''Numeral-Type.bit0''* **shows**
term-of-class.term-of x =
Code-Evaluation.App (*Code-Evaluation.Const* (*STR ''Uint8.uint8.Abs-uint8''*) (*TR* (*STR ''fun''*) [*TR* (*STR ''Word.word''*) [*TR bit0* [*TR bit0* [*TR bit0* [*TR* (*STR ''Numeral-Type.num1''*) []]]]], *TR* (*STR ''Uint8.uint8''*) []])) (*term-of-class.term-of* (*Rep-uint8' x*))
⟨*proof*⟩

lemma *Uin8-code* [*code*]: *Rep-uint8* (*Uint8 i*) = *word-of-int* (*int-of-integer-symbolic i*)
⟨*proof*⟩

code-printing type-constructor *uint8* ⇀
(*SML*) *Word8.word* **and**
(*Haskell*) *Uint8.Word8* **and**
(*Scala*) *Byte*
| **constant** *Uint8* ⇀
(*SML*) *Word8.fromLargeInt* (*IntInf.toLarge* -) **and**
(*Haskell*) (*Prelude.fromInteger* - :: *Uint8.Word8*) **and**
(*Haskell-Quickcheck*) (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint8.Word8*)
and
(*Scala*) -.*byteValue*
| **constant** *0* :: *uint8* ⇀
(*SML*) (*Word8.fromInt 0*) **and**
(*Haskell*) (*0* :: *Uint8.Word8*) **and**
(*Scala*) *0.toByte*
| **constant** *1* :: *uint8* ⇀
(*SML*) (*Word8.fromInt 1*) **and**
(*Haskell*) (*1* :: *Uint8.Word8*) **and**
(*Scala*) *1.toByte*
| **constant** *plus* :: *uint8* ⇒ - ⇒ - ⇀
(*SML*) *Word8.+* ((-), (-)) **and**
(*Haskell*) **infixl** *6* + **and**
(*Scala*) (- +/ -).*toByte*
| **constant** *uminus* :: *uint8* ⇒ - ⇀
(*SML*) *Word8.~* **and**
(*Haskell*) *negate* **and**
(*Scala*) (− -).*toByte*
| **constant** *minus* :: *uint8* ⇒ - ⇀
(*SML*) *Word8.−* ((-), (-)) **and**
(*Haskell*) **infixl** *6* − **and**
(*Scala*) (- −/ -).*toByte*
| **constant** *times* :: *uint8* ⇒ - ⇒ - ⇀

(*SML*) *Word8.** ((-), (-)) **and**
(*Haskell*) **infixl** *7* * **and**
(*Scala*) (- **/* -).*toByte*
| **constant** *HOL.equal* :: *uint8* ⇒ - ⇒ *bool* ⇀
(*SML*) !((- : *Word8.word*) = -) **and**
(*Haskell*) **infix** *4* == **and**
(*Scala*) **infixl** *5* ==
| **class-instance** *uint8* :: *equal* ⇀ (*Haskell*) −
| **constant** *less-eq* :: *uint8* ⇒ - ⇒ *bool* ⇀
(*SML*) *Word8.<=* ((-), (-)) **and**
(*Haskell*) **infix** *4* <= **and**
(*Scala*) *Uint8.less'-eq*
| **constant** *less* :: *uint8* ⇒ - ⇒ *bool* ⇀
(*SML*) *Word8.<* ((-), (-)) **and**
(*Haskell*) **infix** *4* < **and**
(*Scala*) *Uint8.less*
| **constant** *Bit-Operations.not* :: *uint8* ⇒ - ⇀
(*SML*) *Word8.notb* **and**
(*Haskell*) *Data'-Bits.complement* **and**
(*Scala*) -.*unary'-~*.*toByte*
| **constant** *Bit-Operations.and* :: *uint8* ⇒ - ⇀
(*SML*) *Word8.andb* ((-),/ (-)) **and**
(*Haskell*) **infixl** *7 Data-Bits..&.* **and**
(*Scala*) (- & -).*toByte*
| **constant** *Bit-Operations.or* :: *uint8* ⇒ - ⇀
(*SML*) *Word8.orb* ((-),/ (-)) **and**
(*Haskell*) **infixl** *5 Data-Bits..|.* **and**
(*Scala*) (- | -).*toByte*
| **constant** *Bit-Operations.xor* :: *uint8* ⇒ - ⇀
(*SML*) *Word8.xorb* ((-),/ (-)) **and**
(*Haskell*) *Data'-Bits.xor* **and**
(*Scala*) (- ^ -).*toByte*

definition *uint8-divmod* :: *uint8* ⇒ *uint8* ⇒ *uint8* × *uint8* **where**
uint8-divmod x y =
(*if y = 0 then* (*undefined* ((*div*) :: *uint8* ⇒ -) *x* (*0* :: *uint8*), *undefined* ((*mod*) :: *uint8* ⇒ -) *x* (*0* :: *uint8*))
else (*x div y*, *x mod y*))

definition *uint8-div* :: *uint8* ⇒ *uint8* ⇒ *uint8*
where *uint8-div x y* = *fst* (*uint8-divmod x y*)

definition *uint8-mod* :: *uint8* ⇒ *uint8* ⇒ *uint8*
where *uint8-mod x y* = *snd* (*uint8-divmod x y*)

lemma *div-uint8-code* [*code*]: *x div y* = (*if y = 0 then 0 else uint8-div x y*)
including *undefined-transfer* ⟨*proof*⟩

lemma *mod-uint8-code* [*code*]: *x mod y* = (*if y = 0 then x else uint8-mod x y*)

including *undefined-transfer* ⟨*proof*⟩

definition *uint8-sdiv* :: *uint8* ⇒ *uint8* ⇒ *uint8*
where
uint8-sdiv x y =
(*if y* = *0 then undefined* ((*div*) :: *uint8* ⇒ -) *x* (*0* :: *uint8*)
else Abs-uint8 (*Rep-uint8 x sdiv Rep-uint8 y*))

definition *div0-uint8* :: *uint8* ⇒ *uint8*
where [*code del*]: *div0-uint8 x* = *undefined* ((*div*) :: *uint8* ⇒ -) *x* (*0* :: *uint8*)
declare [[*code abort*: *div0-uint8*]]

definition *mod0-uint8* :: *uint8* ⇒ *uint8*
where [*code del*]: *mod0-uint8 x* = *undefined* ((*mod*) :: *uint8* ⇒ -) *x* (*0* :: *uint8*)
declare [[*code abort*: *mod0-uint8*]]

lemma *uint8-divmod-code* [*code*]:
uint8-divmod x y =
(*if 0x80* ≤ *y then if x* < *y then* (*0*, *x*) *else* (*1*, *x* − *y*)
else if y = *0 then* (*div0-uint8 x*, *mod0-uint8 x*)
else let q = *push-bit 1* (*uint8-sdiv* (*drop-bit 1 x*) *y*);
r = *x* − *q* ∗ *y*
in if r ≥ *y then* (*q* + *1*, *r* − *y*) *else* (*q*, *r*))
including *undefined-transfer* ⟨*proof*⟩

lemma *uint8-sdiv-code* [*code*]:
Rep-uint8 (*uint8-sdiv x y*) =
(*if y* = *0 then Rep-uint8* (*undefined* ((*div*) :: *uint8* ⇒ -) *x* (*0* :: *uint8*))
else Rep-uint8 x sdiv Rep-uint8 y)
⟨*proof*⟩

Note that we only need a translation for signed division, but not for the remainder because *uint8-divmod ?x ?y* = (*if 128* ≤ *?y then if ?x* < *?y then* (*0*, *?x*) *else* (*1*, *?x* − *?y*) *else if ?y* = *0 then* (*div0-uint8 ?x*, *mod0-uint8 ?x*) *else let q* = *push-bit 1* (*uint8-sdiv* (*drop-bit 1 ?x*) *?y*); *r* = *?x* − *q* ∗ *?y in if ?y* ≤ *r then* (*q* + *1*, *r* − *?y*) *else* (*q*, *r*)) computes both with division only.

code-printing
constant *uint8-div* ⇀
(*SML*) *Word8.div* ((-), (-)) **and**
(*Haskell*) *Prelude.div*
| **constant** *uint8-mod* ⇀
(*SML*) *Word8.mod* ((-), (-)) **and**
(*Haskell*) *Prelude.mod*
| **constant** *uint8-divmod* ⇀
(*Haskell*) *divmod*
| **constant** *uint8-sdiv* ⇀
(*Scala*) (- *'/* -).*toByte*

definition *uint8-test-bit* :: *uint8* ⇒ *integer* ⇒ *bool*
where [*code del*]:
uint8-test-bit x n =
(*if n* < *0* ∨ *7* < *n then undefined* (*bit* :: *uint8* ⇒ -) *x n*
else bit x (*nat-of-integer n*))

lemma *bit-uint8-code* [*code*]:
bit x n ⟷ *n* < *8* ∧ *uint8-test-bit x* (*integer-of-nat n*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint8-test-bit-code* [*code*]:
uint8-test-bit w n =
(*if n* < *0* ∨ *7* < *n then undefined* (*bit* :: *uint8* ⇒ -) *w n else bit* (*Rep-uint8 w*)
(*nat-of-integer n*))
⟨*proof*⟩

code-printing constant *uint8-test-bit* ⇀
(*SML*) *Uint8.test'-bit* **and**
(*Haskell*) *Data'-Bits.testBitBounded* **and**
(*Scala*) *Uint8.test'-bit* **and**
(*Eval*) (*fn x => fn i => if i < 0 orelse i >= 8 then raise* (*Fail argument to uint8'-test'-bit out of bounds*) *else Uint8.test'-bit x i*)

definition *uint8-set-bit* :: *uint8* ⇒ *integer* ⇒ *bool* ⇒ *uint8*
where [*code del*]:
uint8-set-bit x n b =
(*if n* < *0* ∨ *7* < *n then undefined* (*set-bit* :: *uint8* ⇒ -) *x n b*
else set-bit x (*nat-of-integer n*) *b*)

lemma *set-bit-uint8-code* [*code*]:
set-bit x n b = (*if n* < *8 then uint8-set-bit x* (*integer-of-nat n*) *b else x*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint8-set-bit-code* [*code*]:
Rep-uint8 (*uint8-set-bit w n b*) =
(*if n* < *0* ∨ *7* < *n then Rep-uint8* (*undefined* (*set-bit* :: *uint8* ⇒ -) *w n b*)
else set-bit (*Rep-uint8 w*) (*nat-of-integer n*) *b*)
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint8-set-bit* ⇀
(*SML*) *Uint8.set'-bit* **and**
(*Haskell*) *Data'-Bits.setBitBounded* **and**
(*Scala*) *Uint8.set'-bit* **and**
(*Eval*) (*fn x => fn i => fn b => if i < 0 orelse i >= 8 then raise* (*Fail argument to uint8'-set'-bit out of bounds*) *else Uint8.set'-bit x i b*)

definition *uint8-shiftl* :: *uint8* ⇒ *integer* ⇒ *uint8*

where [*code del*]:
uint8-shiftl x n = (*if n* < *0* ∨ *8* ≤ *n then undefined* (*push-bit* :: *nat* ⇒ *uint8* ⇒ -) *x n else push-bit* (*nat-of-integer n*) *x*)

lemma *shiftl-uint8-code* [*code*]:
push-bit n x = (*if n* < *8 then uint8-shiftl x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint8-shiftl-code* [*code*]:
Rep-uint8 (*uint8-shiftl w n*) =
(*if n* < *0* ∨ *8* ≤ *n then Rep-uint8* (*undefined* (*push-bit* :: *nat* ⇒ *uint8* ⇒ -) *w n*)
else push-bit (*nat-of-integer n*) (*Rep-uint8 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint8-shiftl* ⇀
(*SML*) *Uint8.shiftl* **and**
(*Haskell*) *Data'-Bits.shiftlBounded* **and**
(*Scala*) *Uint8.shiftl* **and**
(*Eval*) (*fn x => fn i => if i* < *0 orelse i* >= *8 then raise* (*Fail argument to uint8'-shiftl out of bounds*) *else Uint8.shiftl x i*)

definition *uint8-shiftr* :: *uint8* ⇒ *integer* ⇒ *uint8*
where [*code del*]:
uint8-shiftr x n = (*if n* < *0* ∨ *8* ≤ *n then undefined* (*drop-bit* :: - ⇒ - ⇒ *uint8*) *x n else drop-bit* (*nat-of-integer n*) *x*)

lemma *shiftr-uint8-code* [*code*]:
drop-bit n x = (*if n* < *8 then uint8-shiftr x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint8-shiftr-code* [*code*]:
Rep-uint8 (*uint8-shiftr w n*) =
(*if n* < *0* ∨ *8* ≤ *n then Rep-uint8* (*undefined* (*drop-bit* :: - ⇒ - ⇒ *uint8*) *w n*)
else drop-bit (*nat-of-integer n*) (*Rep-uint8 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint8-shiftr* ⇀
(*SML*) *Uint8.shiftr* **and**
(*Haskell*) *Data'-Bits.shiftrBounded* **and**
(*Scala*) *Uint8.shiftr* **and**
(*Eval*) (*fn x => fn i => if i* < *0 orelse i* >= *8 then raise* (*Fail argument to uint8'-shiftr out of bounds*) *else Uint8.shiftr x i*)

definition *uint8-sshiftr* :: *uint8* ⇒ *integer* ⇒ *uint8*
where [*code del*]:
uint8-sshiftr x n =
(*if n* < *0* ∨ *8* ≤ *n then undefined signed-drop-bit-uint8 n x else signed-drop-bit-uint8* (*nat-of-integer n*) *x*)

lemma *sshiftr-uint8-code* [*code*]:
signed-drop-bit-uint8 n x =
(*if n* < *8 then uint8-sshiftr x* (*integer-of-nat n*) *else if bit x 7 then* −*1 else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint8-sshiftr-code* [*code*]:
Rep-uint8 (*uint8-sshiftr w n*) =
(*if n* < *0* ∨ *8* ≤ *n then Rep-uint8* (*undefined signed-drop-bit-uint8 n w*)
else signed-drop-bit (*nat-of-integer n*) (*Rep-uint8 w*))
including *undefined-transfer* ⟨*proof*⟩

code-printing constant *uint8-sshiftr* ⇀
(*SML*) *Uint8.shiftr'-signed* **and**
(*Haskell*)
(*Prelude.fromInteger* (*Prelude.toInteger* (*Data'-Bits.shiftrBounded* (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint8.Int8*) -)) :: *Uint8.Word8*) **and**
(*Scala*) *Uint8.shiftr'-signed* **and**
(*Eval*) (*fn x => fn i => if i* < *0 orelse i* >= *8 then raise* (*Fail argument to uint8'-sshiftr out of bounds*) *else Uint8.shiftr'-signed x i*)

context
includes *bit-operations-syntax*
begin

lemma *uint8-msb-test-bit*: *msb x* ⟷ *bit* (*x* :: *uint8*) *7*
⟨*proof*⟩

lemma *msb-uint16-code* [*code*]: *msb x* ⟷ *uint8-test-bit x 7*
⟨*proof*⟩

lemma *uint8-of-int-code* [*code*]:
uint8-of-int i = *Uint8* (*integer-of-int i*)
including *integer.lifting* ⟨*proof*⟩

lemma *int-of-uint8-code* [*code*]:
int-of-uint8 x = *int-of-integer* (*integer-of-uint8 x*)
⟨*proof*⟩

lemma *uint8-of-nat-code* [*code*]:
uint8-of-nat = *uint8-of-int* ∘ *int*
⟨*proof*⟩

lemma *nat-of-uint8-code* [*code*]:
nat-of-uint8 x = *nat-of-integer* (*integer-of-uint8 x*)
⟨*proof*⟩ **including** *integer.lifting* ⟨*proof*⟩

definition *integer-of-uint8-signed* :: *uint8* ⇒ *integer*
where
integer-of-uint8-signed n = (*if bit n 7 then undefined integer-of-uint8 n else inte-*

ger-of-uint8 n)

lemma *integer-of-uint8-signed-code* [*code*]:
integer-of-uint8-signed n =
(*if bit n 7 then undefined integer-of-uint8 n else integer-of-int* (*uint* (*Rep-uint8' n*)))
⟨*proof*⟩

lemma *integer-of-uint8-code* [*code*]:
integer-of-uint8 n =
(*if bit n 7 then integer-of-uint8-signed* (*n AND 0x7F*) *OR 0x80 else integer-of-uint8-signed n*)
⟨*proof*⟩
including *integer.lifting* ⟨*proof*⟩

end

code-printing
constant *integer-of-uint8* ⇀
(*SML*) *IntInf.fromLarge* (*Word8.toLargeInt* -) **and**
(*Haskell*) *Prelude.toInteger*
| **constant** *integer-of-uint8-signed* ⇀
(*Scala*) *BigInt*

9.3 Quickcheck setup

definition *uint8-of-natural* :: *natural* ⇒ *uint8*
where *uint8-of-natural x* ≡ *Uint8* (*integer-of-natural x*)

instantiation *uint8* :: {*random*, *exhaustive*, *full-exhaustive*} **begin**
definition *random-uint8* ≡ *qc-random-cnv uint8-of-natural*
definition *exhaustive-uint8* ≡ *qc-exhaustive-cnv uint8-of-natural*
definition *full-exhaustive-uint8* ≡ *qc-full-exhaustive-cnv uint8-of-natural*
instance ⟨*proof*⟩
end

instantiation *uint8* :: *narrowing* **begin**

interpretation *quickcheck-narrowing-samples*
λi. let x = Uint8 i in (*x*, *0xFF − x*) *0*
Typerep.Typerep (*STR ''Uint8.uint8''*) [] ⟨*proof*⟩

definition *narrowing-uint8 d* = *qc-narrowing-drawn-from* (*narrowing-samples d*) *d*
declare [[*code drop*: *partial-term-of* :: *uint8 itself* ⇒ -]]
lemmas *partial-term-of-uint8* [*code*] = *partial-term-of-code*

instance ⟨*proof*⟩
end

end

Chapter 10

Unsigned words of default size

theory *Uint* **imports**
Code-Target-Word-Base Word-Type-Copies
begin

This theory provides access to words in the target languages of the code generator whose bit width is the default of the target language. To that end, the type *uint* models words of width *dflt-size*, but *dflt-size* is known only to be positive.

Usage restrictions: Default-size words (type *uint*) cannot be used for evaluation, because the results depend on the particular choice of word size in the target language and implementation. Symbolic evaluation has not yet been set up for *uint*.

The default size type

typedecl *dflt-size*

instantiation *dflt-size* :: *typerep* **begin**
definition *typerep-class.typerep* ≡ λ- :: *dflt-size itself*. *Typerep.Typerep* (*STR* ''*Uint.dflt-size*'') []
instance ⟨*proof*⟩
end

consts *dflt-size-aux* :: *nat*
specification (*dflt-size-aux*) *dflt-size-aux-g0*: *dflt-size-aux* > *0*
⟨*proof*⟩
hide-fact *dflt-size-aux-def*

instantiation *dflt-size* :: *len* **begin**
definition *len-of-dflt-size* (- :: *dflt-size itself*) ≡ *dflt-size-aux*
instance ⟨*proof*⟩
end

abbreviation *dflt-size* ≡ *len-of* (*TYPE* (*dflt-size*))

context includes *integer.lifting* **begin**
lift-definition *dflt-size-integer* :: *integer* **is** *int dflt-size* ⟨*proof*⟩
declare *dflt-size-integer-def*[*code del*]
— The code generator will substitute a machine-dependent value for this constant

lemma *dflt-size-by-int*[*code*]: *dflt-size* = *nat-of-integer dflt-size-integer*
⟨*proof*⟩

lemma *dflt-size*[*simp*]:
dflt-size > *0*
dflt-size ≥ *Suc 0*
¬ *dflt-size* < *Suc 0*
⟨*proof*⟩
end

10.1 Type definition and primitive operations

typedef *uint* = ‹*UNIV* :: *dflt-size word set*› ⟨*proof*⟩

global-interpretation *uint*: *word-type-copy Abs-uint Rep-uint*
⟨*proof*⟩

setup-lifting *type-definition-uint*

declare *uint.of-word-of* [*code abstype*]

declare *Quotient-uint* [*transfer-rule*]

instantiation *uint* :: ‹{*comm-ring-1*, *semiring-modulo*, *equal*, *linorder*}›
begin

lift-definition *zero-uint* :: *uint* **is** *0* ⟨*proof*⟩
lift-definition *one-uint* :: *uint* **is** *1* ⟨*proof*⟩
lift-definition *plus-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹(+)› ⟨*proof*⟩
lift-definition *uminus-uint* :: ‹*uint* ⇒ *uint*› **is** *uminus* ⟨*proof*⟩
lift-definition *minus-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹(−)› ⟨*proof*⟩
lift-definition *times-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹(∗)› ⟨*proof*⟩
lift-definition *divide-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹(*div*)› ⟨*proof*⟩
lift-definition *modulo-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹(*mod*)› ⟨*proof*⟩
lift-definition *equal-uint* :: ‹*uint* ⇒ *uint* ⇒ *bool*› **is** ‹*HOL.equal*› ⟨*proof*⟩
lift-definition *less-eq-uint* :: ‹*uint* ⇒ *uint* ⇒ *bool*› **is** ‹(≤)› ⟨*proof*⟩
lift-definition *less-uint* :: ‹*uint* ⇒ *uint* ⇒ *bool*› **is** ‹(<)› ⟨*proof*⟩

global-interpretation *uint*: *word-type-copy-ring Abs-uint Rep-uint*
⟨*proof*⟩

instance ⟨*proof*⟩

end

instantiation *uint* :: *ring-bit-operations*
begin

lift-definition *bit-uint* :: ‹*uint* ⇒ *nat* ⇒ *bool*› **is** *bit* ⟨*proof*⟩
lift-definition *not-uint* :: ‹*uint* ⇒ *uint*› **is** ‹*Bit-Operations.not*› ⟨*proof*⟩
lift-definition *and-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹*Bit-Operations.and*› ⟨*proof*⟩
lift-definition *or-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹*Bit-Operations.or*› ⟨*proof*⟩
lift-definition *xor-uint* :: ‹*uint* ⇒ *uint* ⇒ *uint*› **is** ‹*Bit-Operations.xor*› ⟨*proof*⟩
lift-definition *mask-uint* :: ‹*nat* ⇒ *uint*› **is** *mask* ⟨*proof*⟩
lift-definition *push-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *push-bit* ⟨*proof*⟩
lift-definition *drop-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *drop-bit* ⟨*proof*⟩
lift-definition *signed-drop-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *signed-drop-bit* ⟨*proof*⟩
lift-definition *take-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *take-bit* ⟨*proof*⟩
lift-definition *set-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *Bit-Operations.set-bit* ⟨*proof*⟩
lift-definition *unset-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *unset-bit* ⟨*proof*⟩
lift-definition *flip-bit-uint* :: ‹*nat* ⇒ *uint* ⇒ *uint*› **is** *flip-bit* ⟨*proof*⟩

global-interpretation *uint*: *word-type-copy-bits Abs-uint Rep-uint signed-drop-bit-uint*
⟨*proof*⟩

instance
⟨*proof*⟩

end

lift-definition *uint-of-nat* :: ‹*nat* ⇒ *uint*›
is *word-of-nat* ⟨*proof*⟩

lift-definition *nat-of-uint* :: ‹*uint* ⇒ *nat*›
is *unat* ⟨*proof*⟩

lift-definition *uint-of-int* :: ‹*int* ⇒ *uint*›
is *word-of-int* ⟨*proof*⟩

lift-definition *int-of-uint* :: ‹*uint* ⇒ *int*›
is *uint* ⟨*proof*⟩

context
includes *integer.lifting*
begin

lift-definition *Uint* :: ‹*integer* ⇒ *uint*›
is *word-of-int* ⟨*proof*⟩

lift-definition *integer-of-uint* :: ‹*uint* ⇒ *integer*›

is *uint* ⟨*proof*⟩

end

global-interpretation *uint*: *word-type-copy-more Abs-uint Rep-uint signed-drop-bit-uint uint-of-nat nat-of-uint uint-of-int int-of-uint Uint integer-of-uint* ⟨*proof*⟩

instantiation *uint* :: {*size*, *msb*, *lsb*, *set-bit*, *bit-comprehension*}
begin

lift-definition *size-uint* :: ‹*uint* ⇒ *nat*› **is** *size* ⟨*proof*⟩

lift-definition *msb-uint* :: ‹*uint* ⇒ *bool*› **is** *msb* ⟨*proof*⟩
lift-definition *lsb-uint* :: ‹*uint* ⇒ *bool*› **is** *lsb* ⟨*proof*⟩

Workaround: avoid name space clash by spelling out *lift-definition* explicitly.

definition *set-bit-uint* :: ‹*uint* ⇒ *nat* ⇒ *bool* ⇒ *uint*›
where *set-bit-uint-eq*: ‹*set-bit-uint a n b* = (*if b then Bit-Operations.set-bit else unset-bit*) *n a*›

context
includes *lifting-syntax*
begin

lemma *set-bit-uint-transfer* [*transfer-rule*]:
‹(*cr-uint* ===> (=) ===> (⟷) ===> *cr-uint*) *Generic-set-bit.set-bit Generic-set-bit.set-bit*›
⟨*proof*⟩

end

lift-definition *set-bits-uint* :: ‹(*nat* ⇒ *bool*) ⇒ *uint*› **is** *set-bits* ⟨*proof*⟩
lift-definition *set-bits-aux-uint* :: ‹(*nat* ⇒ *bool*) ⇒ *nat* ⇒ *uint* ⇒ *uint*› **is** *set-bits-aux* ⟨*proof*⟩

global-interpretation *uint*: *word-type-copy-misc Abs-uint Rep-uint signed-drop-bit-uint uint-of-nat nat-of-uint uint-of-int int-of-uint Uint integer-of-uint dflt-size set-bits-aux-uint* ⟨*proof*⟩

instance ⟨*proof*⟩

end

10.2 Code setup

code-printing code-module *Uint* ⇀ (*SML*)
‹
structure Uint : *sig*
val set-bit : *Word.word* −> *IntInf.int* −> *bool* −> *Word.word*

```
  val shiftl : Word.word -> IntInf.int -> Word.word
  val shiftr : Word.word -> IntInf.int -> Word.word
  val shiftr-signed : Word.word -> IntInf.int -> Word.word
  val test-bit : Word.word -> IntInf.int -> bool
end = struct

fun set-bit x n b =
  let val mask = Word.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))
  in if b then Word.orb (x, mask)
     else Word.andb (x, Word.notb mask)
  end

fun shiftl x n =
  Word.<< (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr x n =
  Word.>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun shiftr-signed x n =
  Word.~>> (x, Word.fromLargeInt (IntInf.toLarge n))

fun test-bit x n =
   Word.andb (x, Word.<< (0wx1, Word.fromLargeInt (IntInf.toLarge n))) <>
Word.fromInt 0

end; (* struct Uint *)›
code-reserved SML Uint

code-printing code-module Uint → (Haskell)
 ‹module Uint(Int, Word, dflt-size) where

  import qualified Prelude
  import Data.Int(Int)
  import Data.Word(Word)
  import qualified Data.Bits

  dflt-size :: Prelude.Integer
  dflt-size = Prelude.toInteger (bitSize-aux (0::Word)) where
    bitSize-aux :: (Data.Bits.Bits a, Prelude.Bounded a) => a -> Int
    bitSize-aux = Data.Bits.bitSize›
 and (Haskell-Quickcheck)
 ‹module Uint(Int, Word, dflt-size) where

  import qualified Prelude
  import Data.Int(Int)
  import Data.Word(Word)
  import qualified Data.Bits

  dflt-size :: Prelude.Int
```

```
  dflt-size = bitSize-aux (0::Word) where
    bitSize-aux :: (Data.Bits.Bits a, Prelude.Bounded a) => a -> Int
    bitSize-aux = Data.Bits.bitSize
›
```

code-reserved *Haskell Uint dflt-size*

OCaml and Scala provide only signed bit numbers, so we use these and implement sign-sensitive operations like comparisons manually.

code-printing code-module *Uint* ⇀ (*OCaml*)

```
‹module Uint : sig
  type t = int
  val dflt-size : Z.t
  val less : t -> t -> bool
  val less-eq : t -> t -> bool
  val set-bit : t -> Z.t -> bool -> t
  val shiftl : t -> Z.t -> t
  val shiftr : t -> Z.t -> t
  val shiftr-signed : t -> Z.t -> t
  val test-bit : t -> Z.t -> bool
  val int-mask : int
  val int32-mask : int32
  val int64-mask : int64
end = struct

type t = int

let dflt-size = Z.of-int Sys.int-size;;

(* negative numbers have their highest bit set,
   so they are greater than positive ones *)
let less x y =
  if x<0 then
    y<0 && x<y
  else y < 0 || x < y;;

let less-eq x y =
  if x < 0 then
    y < 0 &&  x <= y
  else y < 0 || x <= y;;

let set-bit x n b =
  let mask = 1 lsl (Z.to-int n)
  in if b then x lor mask
     else x land (lnot mask);;

let shiftl x n = x lsl (Z.to-int n);;

let shiftr x n = x lsr (Z.to-int n);;
```

```
let shiftr-signed x n = x asr (Z.to-int n);;

let test-bit x n = x land (1 lsl (Z.to-int n)) <> 0;;

let int-mask =
  if Sys.int-size < 32 then lnot 0 else 0xFFFFFFFF;;

let int32-mask =
  if Sys.int-size < 32 then Int32.pred (Int32.shift-left Int32.one Sys.int-size)
  else Int32.of-string 0xFFFFFFFF;;

let int64-mask =
  if Sys.int-size < 64 then Int64.pred (Int64.shift-left Int64.one Sys.int-size)
  else Int64.of-string 0xFFFFFFFFFFFFFFFF;;

end;; (*struct Uint*)›
code-reserved OCaml Uint

code-printing code-module Uint ⇀ (Scala)
‹object Uint {
def dflt-size : BigInt = BigInt(32)

def less(x: Int, y: Int) : Boolean =
  if (x < 0) y < 0 && x < y
  else y < 0 || x < y

def less-eq(x: Int, y: Int) : Boolean =
  if (x < 0) y < 0 && x <= y
  else y < 0 || x <= y

def set-bit(x: Int, n: BigInt, b: Boolean) : Int =
  if (b)
    x | (1 << n.intValue)
  else
    x & (1 << n.intValue).unary-~

def shiftl(x: Int, n: BigInt) : Int = x << n.intValue

def shiftr(x: Int, n: BigInt) : Int = x >>> n.intValue

def shiftr-signed(x: Int, n: BigInt) : Int = x >> n.intValue

def test-bit(x: Int, n: BigInt) : Boolean =
  (x & (1 << n.intValue)) != 0

} /* object Uint */›
code-reserved Scala Uint
```

OCaml's conversion from Big_int to int demands that the value fits into a

signed integer. The following justifies the implementation.

context
includes *integer.lifting bit-operations-syntax*
begin

definition *wivs-mask* :: *int* **where** *wivs-mask* = *2^ dflt-size − 1*
lift-definition *wivs-mask-integer* :: *integer* **is** *wivs-mask* ⟨*proof*⟩
lemma [*code*]: *wivs-mask-integer* = *2 ^ dflt-size − 1*
⟨*proof*⟩

definition *wivs-shift* :: *int* **where** *wivs-shift* = *2 ^ dflt-size*
lift-definition *wivs-shift-integer* :: *integer* **is** *wivs-shift* ⟨*proof*⟩
lemma [*code*]: *wivs-shift-integer* = *2 ^ dflt-size*
⟨*proof*⟩

definition *wivs-index* :: *nat* **where** *wivs-index* == *dflt-size − 1*
lift-definition *wivs-index-integer* :: *integer* **is** *int wivs-index*⟨*proof*⟩
lemma *wivs-index-integer-code*[*code*]: *wivs-index-integer* = *dflt-size-integer − 1*
⟨*proof*⟩

definition *wivs-overflow* :: *int* **where** *wivs-overflow* == *2^ (dflt-size − 1)*
lift-definition *wivs-overflow-integer* :: *integer* **is** *wivs-overflow* ⟨*proof*⟩
lemma [*code*]: *wivs-overflow-integer* = *2 ^ (dflt-size − 1)*
⟨*proof*⟩

definition *wivs-least* :: *int* **where** *wivs-least* == *− wivs-overflow*
lift-definition *wivs-least-integer* :: *integer* **is** *wivs-least* ⟨*proof*⟩
lemma [*code*]: *wivs-least-integer* = *− (2 ^ (dflt-size − 1))*
⟨*proof*⟩

definition *Uint-signed* :: *integer* ⇒ *uint* **where**
Uint-signed i = (*if i < wivs-least-integer ∨ wivs-overflow-integer ≤ i then undefined Uint i else Uint i*)

lemma *Uint-code* [*code*]:
Uint i =
(*let i′* = *i AND wivs-mask-integer in*
if bit i′ wivs-index then Uint-signed (*i′ − wivs-shift-integer*) *else Uint-signed i′*)
including *undefined-transfer*
⟨*proof*⟩

lemma *Uint-signed-code* [*code*]:
Rep-uint (*Uint-signed i*) =
(*if i < wivs-least-integer ∨ i ≥ wivs-overflow-integer then Rep-uint* (*undefined Uint i*) *else word-of-int* (*int-of-integer-symbolic i*))
⟨*proof*⟩
end

Avoid *Abs-uint* in generated code, use *Rep-uint′* instead. The symbolic im-

plementations for code_simp use *Rep-uint*.

The new destructor *Rep-uint′* is executable. As the simplifier is given the [code abstract] equations literally, we cannot implement *Rep-uint* directly, because that makes code_simp loop.

If code generation raises Match, some equation probably contains *Rep-uint* ([code abstract] equations for *uint* may use *Rep-uint* because these instances will be folded away.)

definition *Rep-uint′* **where** [*simp*]: *Rep-uint′* = *Rep-uint*

lemma *Rep-uint′-code* [*code*]: *Rep-uint′ x* = (*BITS n. bit x n*)
⟨*proof*⟩

lift-definition *Abs-uint′* :: *dflt-size word* ⇒ *uint* **is** λ*x* :: *dflt-size word. x* ⟨*proof*⟩

lemma *Abs-uint′-code* [*code*]:
Abs-uint′ x = *Uint* (*integer-of-int* (*uint x*))
including *integer.lifting* ⟨*proof*⟩

declare [[*code drop*: *term-of-class.term-of* :: *uint* ⇒ -]]

lemma *term-of-uint-code* [*code*]:
defines *TR* ≡ *typerep.Typerep* **and** *bit0* ≡ *STR ''Numeral-Type.bit0''*
shows
term-of-class.term-of x =
Code-Evaluation.App (*Code-Evaluation.Const* (*STR ''Uint.uint.Abs-uint''*) (*TR* (*STR ''fun''*) [*TR* (*STR ''Word.word''*) [*TR* (*STR ''Uint.dflt-size''*) []], *TR* (*STR ''Uint.uint''*) []]))
(*term-of-class.term-of* (*Rep-uint′ x*))
⟨*proof*⟩

Important: We must prevent the reflection oracle (eval-tac) to use our machine-dependent type.

code-printing
type-constructor *uint* ⇀
(*SML*) *Word.word* **and**
(*Haskell*) *Uint.Word* **and**
(*OCaml*) *Uint.t* **and**
(*Scala*) *Int* **and**
(*Eval*) ∗∗∗ *Error*: *Machine dependent type* ∗∗∗ **and**
(*Quickcheck*) *Word.word*
| **constant** *dflt-size-integer* ⇀
(*SML*) (*IntInf.fromLarge* (*Int.toLarge Word.wordSize*)) **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Word.wordSize* **and**
(*Haskell*) *Uint.dflt′-size* **and**
(*OCaml*) *Uint.dflt′-size* **and**
(*Scala*) *Uint.dflt′-size*
| **constant** *Uint* ⇀

```
  (SML) Word.fromLargeInt (IntInf.toLarge -) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.fromInt and
  (Haskell) (Prelude.fromInteger - :: Uint.Word) and
  (Haskell-Quickcheck) (Prelude.fromInteger (Prelude.toInteger -) :: Uint.Word)
and
  (Scala) -.intValue
| constant Uint-signed →
  (OCaml) Z.to'-int
| constant 0 :: uint →
  (SML) (Word.fromInt 0) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) (Word.fromInt 0) and
  (Haskell) (0 :: Uint.Word) and
  (OCaml) 0 and
  (Scala) 0
| constant 1 :: uint →
  (SML) (Word.fromInt 1) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) (Word.fromInt 1) and
  (Haskell) (1 :: Uint.Word) and
  (OCaml) 1 and
  (Scala) 1
| constant plus :: uint ⇒ - →
  (SML) Word.+ ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.+ ((-), (-)) and
  (Haskell) infixl 6 + and
  (OCaml) Pervasives.(+) and
  (Scala) infixl 7 +
| constant uminus :: uint ⇒ - →
  (SML) Word.~ and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.~ and
  (Haskell) negate and
  (OCaml) Pervasives.(~−) and
  (Scala) !(− -)
| constant minus :: uint ⇒ - →
  (SML) Word.− ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.− ((-), (-)) and
  (Haskell) infixl 6 − and
  (OCaml) Pervasives.(−) and
  (Scala) infixl 7 −
| constant times :: uint ⇒ - ⇒ - →
  (SML) Word.* ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.* ((-), (-)) and
  (Haskell) infixl 7 * and
```

```
  (OCaml) Pervasives.( * ) and
  (Scala) infixl 8 *
| constant HOL.equal :: uint ⇒ - ⇒ bool ⇀
  (SML) !((- : Word.word) = -) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) !((- : Word.word) = -) and
  (Haskell) infix 4 == and
  (OCaml) (Pervasives.(=):Uint.t -> Uint.t -> bool) and
  (Scala) infixl 5 ==
| class-instance uint :: equal ⇀
  (Haskell) -
| constant less-eq :: uint ⇒ - ⇒ bool ⇀
  (SML) Word.<= ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.<= ((-), (-)) and
  (Haskell) infix 4 <= and
  (OCaml) Uint.less'-eq and
  (Scala) Uint.less'-eq
| constant less :: uint ⇒ - ⇒ bool ⇀
  (SML) Word.< ((-), (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.< ((-), (-)) and
  (Haskell) infix 4 < and
  (OCaml) Uint.less and
  (Scala) Uint.less
| constant Bit-Operations.not :: uint ⇒ - ⇀
  (SML) Word.notb and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.notb and
  (Haskell) Data'-Bits.complement and
  (OCaml) Pervasives.lnot and
  (Scala) -.unary'-~
| constant Bit-Operations.and :: uint ⇒ - ⇀
  (SML) Word.andb ((-),/ (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.andb ((-),/ (-)) and
  (Haskell) infixl 7 Data-Bits..&. and
  (OCaml) Pervasives.(land) and
  (Scala) infixl 3 &
| constant Bit-Operations.or :: uint ⇒ - ⇀
  (SML) Word.orb ((-),/ (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
  (Quickcheck) Word.orb ((-),/ (-)) and
  (Haskell) infixl 5 Data-Bits..|. and
  (OCaml) Pervasives.(lor) and
  (Scala) infixl 1 |
| constant Bit-Operations.xor :: uint ⇒ - ⇀
  (SML) Word.xorb ((-),/ (-)) and
  (Eval) (raise (Fail Machine dependent code)) and
```

(*Quickcheck*) *Word.xorb* ((-),/ (-)) **and**
(*Haskell*) *Data'-Bits.xor* **and**
(*OCaml*) *Pervasives.*(*lxor*) **and**
(*Scala*) **infixl** *2* ^

definition *uint-divmod* :: *uint* ⇒ *uint* ⇒ *uint* × *uint* **where**
uint-divmod x y =
(*if y* = *0 then* (*undefined* ((*div*) :: *uint* ⇒ -) *x* (*0* :: *uint*), *undefined* ((*mod*) :: *uint* ⇒ -) *x* (*0* :: *uint*))
else (*x div y*, *x mod y*))

definition *uint-div* :: *uint* ⇒ *uint* ⇒ *uint*
where *uint-div x y* = *fst* (*uint-divmod x y*)

definition *uint-mod* :: *uint* ⇒ *uint* ⇒ *uint*
where *uint-mod x y* = *snd* (*uint-divmod x y*)

lemma *div-uint-code* [*code*]: *x div y* = (*if y* = *0 then 0 else uint-div x y*)
including *undefined-transfer* ⟨*proof*⟩

lemma *mod-uint-code* [*code*]: *x mod y* = (*if y* = *0 then x else uint-mod x y*)
including *undefined-transfer* ⟨*proof*⟩

definition *uint-sdiv* :: *uint* ⇒ *uint* ⇒ *uint*
where [*code del*]:
uint-sdiv x y =
(*if y* = *0 then undefined* ((*div*) :: *uint* ⇒ -) *x* (*0* :: *uint*)
else Abs-uint (*Rep-uint x sdiv Rep-uint y*))

definition *div0-uint* :: *uint* ⇒ *uint*
where [*code del*]: *div0-uint x* = *undefined* ((*div*) :: *uint* ⇒ -) *x* (*0* :: *uint*)
declare [[*code abort*: *div0-uint*]]

definition *mod0-uint* :: *uint* ⇒ *uint*
where [*code del*]: *mod0-uint x* = *undefined* ((*mod*) :: *uint* ⇒ -) *x* (*0* :: *uint*)
declare [[*code abort*: *mod0-uint*]]

definition *wivs-overflow-uint* :: *uint*
where *wivs-overflow-uint* ≡ *push-bit* (*dflt-size* − *1*) *1*

lemma *Rep-uint-wivs-overflow-uint-eq*:
‹*Rep-uint wivs-overflow-uint* = *2* ^ (*dflt-size* − *Suc 0*)›
⟨*proof*⟩

lemma *wivs-overflow-uint-greater-eq-0*:
‹*wivs-overflow-uint* > *0*›
⟨*proof*⟩

lemma *uint-divmod-code* [*code*]:

uint-divmod x y =
(if wivs-overflow-uint ≤ y then if x < y then (0, x) else (1, x − y)
else if y = 0 then (div0-uint x, mod0-uint x)
else let q = push-bit 1 (uint-sdiv (drop-bit 1 x) y);
r = x − q ∗ y
in if r ≥ y then (q + 1, r − y) else (q, r))
⟨*proof*⟩
including *undefined-transfer*
⟨*proof*⟩

lemma *uint-sdiv-code* [*code*]:
Rep-uint (uint-sdiv x y) =
(if y = 0 then Rep-uint (undefined ((div) :: uint ⇒ -) x (0 :: uint))
else Rep-uint x sdiv Rep-uint y)
⟨*proof*⟩

Note that we only need a translation for signed division, but not for the remainder because *uint-divmod ?x ?y = (if wivs-overflow-uint ≤ ?y then if ?x < ?y then (0, ?x) else (1, ?x − ?y) else if ?y = 0 then (div0-uint ?x, mod0-uint ?x) else let q = push-bit 1 (uint-sdiv (drop-bit 1 ?x) ?y); r = ?x − q ∗ ?y in if ?y ≤ r then (q + 1, r − ?y) else (q, r))* computes both with division only.

code-printing
constant *uint-div* ⇀
(*SML*) *Word.div* ((-), (-)) **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Word.div* ((-), (-)) **and**
(*Haskell*) *Prelude.div*
| **constant** *uint-mod* ⇀
(*SML*) *Word.mod* ((-), (-)) **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Word.mod* ((-), (-)) **and**
(*Haskell*) *Prelude.mod*
| **constant** *uint-divmod* ⇀
(*Haskell*) *divmod*
| **constant** *uint-sdiv* ⇀
(*OCaml*) *Pervasives.*(*'/*) **and**
(*Scala*) - *'/* -

definition *uint-test-bit* :: *uint* ⇒ *integer* ⇒ *bool*
where [*code del*]:
uint-test-bit x n =
(if n < 0 ∨ dflt-size-integer ≤ n then undefined (bit :: uint ⇒ -) x n
else bit x (nat-of-integer n))

lemma *test-bit-uint-code* [*code*]:
bit x n ⟷ n < dflt-size ∧ uint-test-bit x (integer-of-nat n)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint-test-bit-code* [*code*]:
uint-test-bit w n =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then undefined* (*bit* :: *uint* ⇒ -) *w n else bit* (*Rep-uint w*) (*nat-of-integer n*))
⟨*proof*⟩

code-printing constant *uint-test-bit* ⇀
(*SML*) *Uint.test'-bit* **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Uint.test'-bit* **and**
(*Haskell*) *Data'-Bits.testBitBounded* **and**
(*OCaml*) *Uint.test'-bit* **and**
(*Scala*) *Uint.test'-bit*

definition *uint-set-bit* :: *uint* ⇒ *integer* ⇒ *bool* ⇒ *uint*
where [*code del*]:
uint-set-bit x n b =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then undefined* (*set-bit* :: *uint* ⇒ -) *x n b*
else set-bit x (*nat-of-integer n*) *b*)

lemma *set-bit-uint-code* [*code*]:
set-bit x n b = (*if n* < *dflt-size then uint-set-bit x* (*integer-of-nat n*) *b else x*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint-set-bit-code* [*code*]:
Rep-uint (*uint-set-bit w n b*) =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then Rep-uint* (*undefined* (*set-bit* :: *uint* ⇒ -) *w n b*)
else set-bit (*Rep-uint w*) (*nat-of-integer n*) *b*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

code-printing constant *uint-set-bit* ⇀
(*SML*) *Uint.set'-bit* **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Uint.set'-bit* **and**
(*Haskell*) *Data'-Bits.setBitBounded* **and**
(*OCaml*) *Uint.set'-bit* **and**
(*Scala*) *Uint.set'-bit*

definition *uint-shiftl* :: *uint* ⇒ *integer* ⇒ *uint*
where [*code del*]:
uint-shiftl x n = (*if n* < *0* ∨ *dflt-size-integer* ≤ *n then undefined* (*push-bit* :: *nat* ⇒ *uint* ⇒ -) *x n else push-bit* (*nat-of-integer n*) *x*)

lemma *shiftl-uint-code* [*code*]: *push-bit n x* = (*if n* < *dflt-size then uint-shiftl x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint-shiftl-code* [*code*]:

Rep-uint (*uint-shiftl w n*) =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then Rep-uint* (*undefined* (*push-bit* :: *nat* ⇒ *uint* ⇒ -) *w n*) *else push-bit* (*nat-of-integer n*) (*Rep-uint w*))
including *undefined-transfer integer.lifting* ⟨*proof*⟩

code-printing constant *uint-shiftl* ⇀
(*SML*) *Uint.shiftl* **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Uint.shiftl* **and**
(*Haskell*) *Data′-Bits.shiftlBounded* **and**
(*OCaml*) *Uint.shiftl* **and**
(*Scala*) *Uint.shiftl*

definition *uint-shiftr* :: *uint* ⇒ *integer* ⇒ *uint*
where [*code del*]:
uint-shiftr x n = (*if n* < *0* ∨ *dflt-size-integer* ≤ *n then undefined* (*drop-bit* :: *nat* ⇒ *uint* ⇒ -) *x n else drop-bit* (*nat-of-integer n*) *x*)

lemma *shiftr-uint-code* [*code*]: *drop-bit n x* = (*if n* < *dflt-size then uint-shiftr x* (*integer-of-nat n*) *else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint-shiftr-code* [*code*]:
Rep-uint (*uint-shiftr w n*) =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then Rep-uint* (*undefined* (*drop-bit* :: *nat* ⇒ *uint* ⇒ -) *w n*) *else drop-bit* (*nat-of-integer n*) (*Rep-uint w*))
including *undefined-transfer integer.lifting* ⟨*proof*⟩

code-printing constant *uint-shiftr* ⇀
(*SML*) *Uint.shiftr* **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Uint.shiftr* **and**
(*Haskell*) *Data′-Bits.shiftrBounded* **and**
(*OCaml*) *Uint.shiftr* **and**
(*Scala*) *Uint.shiftr*

definition *uint-sshiftr* :: *uint* ⇒ *integer* ⇒ *uint*
where [*code del*]:
uint-sshiftr x n =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then undefined signed-drop-bit-uint n x else signed-drop-bit-uint* (*nat-of-integer n*) *x*)

lemma *sshiftr-uint-code* [*code*]:
signed-drop-bit-uint n x =
(*if n* < *dflt-size then uint-sshiftr x* (*integer-of-nat n*) *else*
if bit x wivs-index then −*1 else 0*)
including *undefined-transfer integer.lifting* ⟨*proof*⟩

lemma *uint-sshiftr-code* [*code*]:

Rep-uint (*uint-sshiftr w n*) =
(*if n* < *0* ∨ *dflt-size-integer* ≤ *n then Rep-uint* (*undefined signed-drop-bit-uint n w*) *else signed-drop-bit* (*nat-of-integer n*) (*Rep-uint w*))
including *undefined-transfer integer.lifting* ⟨*proof*⟩

code-printing constant *uint-sshiftr* ⟶
(*SML*) *Uint.shiftr'-signed* **and**
(*Eval*) (*raise* (*Fail Machine dependent code*)) **and**
(*Quickcheck*) *Uint.shiftr'-signed* **and**
(*Haskell*)
(*Prelude.fromInteger* (*Prelude.toInteger* (*Data'-Bits.shiftrBounded* (*Prelude.fromInteger* (*Prelude.toInteger* -) :: *Uint.Int*) -)) :: *Uint.Word*) **and**
(*OCaml*) *Uint.shiftr'-signed* **and**
(*Scala*) *Uint.shiftr'-signed*

lemma *uint-msb-test-bit*: *msb x* ⟷ *bit* (*x* :: *uint*) *wivs-index*
⟨*proof*⟩

lemma *msb-uint-code* [*code*]: *msb x* ⟷ *uint-test-bit x wivs-index-integer*
⟨*proof*⟩

lemma *uint-of-int-code* [*code*]: *uint-of-int i* = (*BITS n*. *bit i n*)
⟨*proof*⟩

10.3 Quickcheck setup

definition *uint-of-natural* :: *natural* ⇒ *uint*
where *uint-of-natural x* ≡ *Uint* (*integer-of-natural x*)

instantiation *uint* :: {*random*, *exhaustive*, *full-exhaustive*} **begin**
definition *random-uint* ≡ *qc-random-cnv uint-of-natural*
definition *exhaustive-uint* ≡ *qc-exhaustive-cnv uint-of-natural*
definition *full-exhaustive-uint* ≡ *qc-full-exhaustive-cnv uint-of-natural*
instance ⟨*proof*⟩
end

instantiation *uint* :: *narrowing* **begin**

interpretation *quickcheck-narrowing-samples*
λi. (*Uint i*, *Uint* (− *i*)) *0*
Typerep.Typerep (*STR ''Uint.uint''*) [] ⟨*proof*⟩

definition *narrowing-uint d* = *qc-narrowing-drawn-from* (*narrowing-samples d*) *d*
declare [[*code drop*: *partial-term-of* :: *uint itself* ⇒ -]]
lemmas *partial-term-of-uint* [*code*] = *partial-term-of-code*

instance ⟨*proof*⟩
end

end

Chapter 11

Conversions between unsigned words and between char

theory *Native-Cast*
imports
Uint8
Uint16
Uint32
Uint64
begin

Auxiliary stuff

lemma *integer-of-char-char-of-integer* [*simp*]:
integer-of-char (*char-of-integer x*) = *x mod 256*
⟨*proof*⟩

lemma *char-of-integer-integer-of-char* [*simp*]:
char-of-integer (*integer-of-char x*) = *x*
⟨*proof*⟩

lemma *int-lt-numeral* [*simp*]: *int x* < *numeral n* ⟷ *x* < *numeral n*
⟨*proof*⟩

lemma *int-of-integer-ge-0*: *0* ≤ *int-of-integer x* ⟷ *0* ≤ *x*
including *integer.lifting* ⟨*proof*⟩

lemma *integer-of-char-ge-0* [*simp*]: *0* ≤ *integer-of-char x*
including *integer.lifting* ⟨*proof*⟩

11.1 Conversion between native words

lift-definition *uint8-of-uint16* :: *uint16* ⇒ *uint8* **is** *ucast* ⟨*proof*⟩

lift-definition *uint8-of-uint32* :: *uint32* ⇒ *uint8* **is** *ucast* ⟨*proof*⟩
lift-definition *uint8-of-uint64* :: *uint64* ⇒ *uint8* **is** *ucast* ⟨*proof*⟩

lift-definition *uint16-of-uint8* :: *uint8* ⇒ *uint16* **is** *ucast* ⟨*proof*⟩
lift-definition *uint16-of-uint32* :: *uint32* ⇒ *uint16* **is** *ucast* ⟨*proof*⟩
lift-definition *uint16-of-uint64* :: *uint64* ⇒ *uint16* **is** *ucast* ⟨*proof*⟩

lift-definition *uint32-of-uint8* :: *uint8* ⇒ *uint32* **is** *ucast* ⟨*proof*⟩
lift-definition *uint32-of-uint16* :: *uint16* ⇒ *uint32* **is** *ucast* ⟨*proof*⟩
lift-definition *uint32-of-uint64* :: *uint64* ⇒ *uint32* **is** *ucast* ⟨*proof*⟩

lift-definition *uint64-of-uint8* :: *uint8* ⇒ *uint64* **is** *ucast* ⟨*proof*⟩
lift-definition *uint64-of-uint16* :: *uint16* ⇒ *uint64* **is** *ucast* ⟨*proof*⟩
lift-definition *uint64-of-uint32* :: *uint32* ⇒ *uint64* **is** *ucast* ⟨*proof*⟩

context
begin

qualified definition *mask* :: *integer*
  where ‹*mask* = (*0xFFFFFFFF* :: *integer*)›

end

code-printing
  constant *uint8-of-uint16* ⇀
  (*SML-word*) *Word8.fromLarge* (*Word16.toLarge* -) **and**
  (*Haskell*) (*Prelude.fromIntegral* - :: *Uint8.Word8*) **and**
  (*Scala*) -.*toByte*
| **constant** *uint8-of-uint32* ⇀
  (*SML*) *Word8.fromLarge* (*Word32.toLarge* -) **and**
  (*Haskell*) (*Prelude.fromIntegral* - :: *Uint8.Word8*) **and**
  (*Scala*) -.*toByte*
| **constant** *uint8-of-uint64* ⇀
  (*SML*) *Word8.fromLarge* (*Uint64.toLarge* -) **and**
  (*Haskell*) (*Prelude.fromIntegral* - :: *Uint8.Word8*) **and**
  (*Scala*) -.*toByte*
| **constant** *uint16-of-uint8* ⇀
  (*SML-word*) *Word16.fromLarge* (*Word8.toLarge* -) **and**
  (*Haskell*) (*Prelude.fromIntegral* - :: *Uint16.Word16*) **and**
  (*Scala*) ((-).*toInt* & *0xFF*).*toChar*
| **constant** *uint16-of-uint32* ⇀
  (*SML-word*) *Word16.fromLarge* (*Word32.toLarge* -) **and**
  (*Haskell*) (*Prelude.fromIntegral* - :: *Uint16.Word16*) **and**
  (*Scala*) -.*toChar*
| **constant** *uint16-of-uint64* ⇀
  (*SML-word*) *Word16.fromLarge* (*Uint64.toLarge* -) **and**
  (*Haskell*) (*Prelude.fromIntegral* - :: *Uint16.Word16*) **and**
  (*Scala*) -.*toChar*
| **constant** *uint32-of-uint8* ⇀

```
  (SML) Word32.fromLarge (Word8.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
  (Scala) ((-).toInt & 0xFF)
| constant uint32-of-uint16 ⇀
  (SML-word) Word32.fromLarge (Word16.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
  (Scala) (-).toInt
| constant uint32-of-uint64 ⇀
  (SML-word) Word32.fromLarge (Uint64.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
  (Scala) (-).toInt and
  (OCaml) Int64.to'-int32
| constant uint64-of-uint8 ⇀
  (SML-word) Word64.fromLarge (Word8.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
  (Scala) ((-).toLong & 0xFF)
| constant uint64-of-uint16 ⇀
  (SML-word) Word64.fromLarge (Word16.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
  (Scala) -.toLong
| constant uint64-of-uint32 ⇀
  (SML-word) Word64.fromLarge (Word32.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
  (Scala) ((-).toLong & 0xFFFFFFFFL) and
  (OCaml) Int64.logand (Int64.of'-int32 -) (Int64.of'-string 4294967295)
```

Use *Abs-uint8'* etc. instead of *Rep-uint8* in code equations for conversion functions to avoid exceptions during code generation when the target language provides only some of the uint types.

lemma *uint8-of-uint16-code* [*code*]:
  uint8-of-uint16 x = Abs-uint8' (ucast (Rep-uint16' x))
⟨*proof*⟩

lemma *uint8-of-uint32-code* [*code*]:
  uint8-of-uint32 x = Abs-uint8' (ucast (Rep-uint32' x))
⟨*proof*⟩

lemma *uint8-of-uint64-code* [*code*]:
  uint8-of-uint64 x = Abs-uint8' (ucast (Rep-uint64' x))
⟨*proof*⟩

lemma *uint16-of-uint8-code* [*code*]:
  uint16-of-uint8 x = Abs-uint16' (ucast (Rep-uint8' x))
⟨*proof*⟩

lemma *uint16-of-uint32-code* [*code*]:
  uint16-of-uint32 x = Abs-uint16' (ucast (Rep-uint32' x))
⟨*proof*⟩

lemma *uint16-of-uint64-code* [*code*]:
uint16-of-uint64 x = *Abs-uint16'* (*ucast* (*Rep-uint64' x*))
⟨*proof*⟩

lemma *uint32-of-uint8-code* [*code*]:
uint32-of-uint8 x = *Abs-uint32'* (*ucast* (*Rep-uint8' x*))
⟨*proof*⟩

lemma *uint32-of-uint16-code* [*code*]:
uint32-of-uint16 x = *Abs-uint32'* (*ucast* (*Rep-uint16' x*))
⟨*proof*⟩

lemma *uint32-of-uint64-code* [*code*]:
uint32-of-uint64 x = *Abs-uint32'* (*ucast* (*Rep-uint64' x*))
⟨*proof*⟩

lemma *uint64-of-uint8-code* [*code*]:
uint64-of-uint8 x = *Abs-uint64'* (*ucast* (*Rep-uint8' x*))
⟨*proof*⟩

lemma *uint64-of-uint16-code* [*code*]:
uint64-of-uint16 x = *Abs-uint64'* (*ucast* (*Rep-uint16' x*))
⟨*proof*⟩

lemma *uint64-of-uint32-code* [*code*]:
uint64-of-uint32 x = *Abs-uint64'* (*ucast* (*Rep-uint32' x*))
⟨*proof*⟩

end

theory *Native-Cast-Uint* **imports**
Native-Cast
Uint
begin

lift-definition *uint-of-uint8* :: *uint8* ⇒ *uint* **is** *ucast* ⟨*proof*⟩
lift-definition *uint-of-uint16* :: *uint16* ⇒ *uint* **is** *ucast* ⟨*proof*⟩
lift-definition *uint-of-uint32* :: *uint32* ⇒ *uint* **is** *ucast* ⟨*proof*⟩
lift-definition *uint-of-uint64* :: *uint64* ⇒ *uint* **is** *ucast* ⟨*proof*⟩

lift-definition *uint8-of-uint* :: *uint* ⇒ *uint8* **is** *ucast* ⟨*proof*⟩
lift-definition *uint16-of-uint* :: *uint* ⇒ *uint16* **is** *ucast* ⟨*proof*⟩
lift-definition *uint32-of-uint* :: *uint* ⇒ *uint32* **is** *ucast* ⟨*proof*⟩
lift-definition *uint64-of-uint* :: *uint* ⇒ *uint64* **is** *ucast* ⟨*proof*⟩

code-printing
constant *uint-of-uint8* ⇀
(*SML*) *Word.fromLarge* (*Word8.toLarge* -) **and**

```
  (Haskell) (Prelude.fromIntegral - :: Uint.Word) and
  (Scala) ((-).toInt & 0xFF)
| constant uint-of-uint16 ⇀
  (SML-word) Word.fromLarge (Word16.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint.Word) and
  (Scala) (-).toInt
| constant uint-of-uint32 ⇀
  (SML) Word.fromLarge (Word32.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint.Word) and
  (Scala) - and
  (OCaml) (Int32.to'-int -) land Uint.int'-mask
| constant uint-of-uint64 ⇀
  (SML) Word.fromLarge (Uint64.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint.Word) and
  (Scala) (-).toInt and
  (OCaml) Int64.to'-int
| constant uint8-of-uint ⇀
  (SML) Word8.fromLarge (Word.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint8.Word8) and
  (Scala) (-).toByte
| constant uint16-of-uint ⇀
  (SML-word) Word16.fromLarge (Word.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint16.Word16) and
  (Scala) (-).toChar
| constant uint32-of-uint ⇀
  (SML) Word32.fromLarge (Word.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint32.Word32) and
  (Scala) - and
  (OCaml) Int32.logand (Int32.of'-int -) Uint.int32'-mask
| constant uint64-of-uint ⇀
  (SML) Uint64.fromLarge (Word.toLarge -) and
  (Haskell) (Prelude.fromIntegral - :: Uint64.Word64) and
  (Scala) ((-).toLong & 0xFFFFFFFFL) and
  (OCaml) Int64.logand (Int64.of'-int -) Uint.int64'-mask

lemma uint8-of-uint-code [code]:
  uint8-of-uint x = Abs-uint8' (ucast (Rep-uint' x))
  ⟨proof⟩

lemma uint16-of-uint-code [code]:
  uint16-of-uint x = Abs-uint16' (ucast (Rep-uint' x))
  ⟨proof⟩

lemma uint32-of-uint-code [code]:
  uint32-of-uint x = Abs-uint32' (ucast (Rep-uint' x))
  ⟨proof⟩

lemma uint64-of-uint-code [code]:
  uint64-of-uint x = Abs-uint64' (ucast (Rep-uint' x))
```

⟨*proof*⟩

lemma *uint-of-uint8-code* [*code*]:
uint-of-uint8 x = *Abs-uint'* (*ucast* (*Rep-uint8' x*))
⟨*proof*⟩

lemma *uint-of-uint16-code* [*code*]:
uint-of-uint16 x = *Abs-uint'* (*ucast* (*Rep-uint16' x*))
⟨*proof*⟩

lemma *uint-of-uint32-code* [*code*]:
uint-of-uint32 x = *Abs-uint'* (*ucast* (*Rep-uint32' x*))
⟨*proof*⟩

lemma *uint-of-uint64-code* [*code*]:
uint-of-uint64 x = *Abs-uint'* (*ucast* (*Rep-uint64' x*))
⟨*proof*⟩

end

11.2 Compatibility with Imperative/HOL

theory *Native-Word-Imperative-HOL* **imports**
Code-Target-Word-Base
HOL−Imperative-HOL.Heap-Monad
begin

We add a code target that combines the translations for native words that are by default not supported by all PolyML versions with the adaptations for Imperative_HOL.

⟨*ML*⟩

end

Chapter 12

Test cases

theory *Native-Word-Test* **imports**
Uint64 Uint32 Uint16 Uint8 Uint Native-Cast-Uint
HOL−Library.Code-Test Word-Lib.Bit-Shifts-Infix-Syntax
begin

export-code
nat-of-uint8 uint8-of-nat
nat-of-uint16 uint16-of-nat
nat-of-uint32 uint32-of-nat
nat-of-uint64 uint64-of-nat
nat-of-uint uint-of-nat
in *SML*

12.1 Tests for *uint32*

context
includes *bit-operations-syntax*
begin

abbreviation (*input*) *sshiftr-uint32* (**infixl** *>>>* *55*)
where ‹*w >>> n* ≡ *signed-drop-bit-uint32 n w*›

definition *test-uint32* **where**
test-uint32 ⟷
(([*0x100000001*, *−1*, *−4294967291*, *0xFFFFFFFF*, *0x12345678*
, *0x5A AND 0x36*
, *0x5A OR 0x36*
, *0x5A XOR 0x36*
, *NOT 0x5A*
, *5 + 6*, *−5 + 6*, *−6 + 5*, *−5 + (− 6)*, *0xFFFFFFFFF + 1*
, *5 − 3*, *3 − 5*
, *5 * 3*, *−5 * 3*, *−5 * −4*, *0x12345678 * 0x87654321*
, *5 div 3*, *−5 div 3*, *−5 div −3*, *5 div −3*
, *5 mod 3*, *−5 mod 3*, *−5 mod −3*, *5 mod −3*

, *set-bit 5 4 True*, *set-bit* (− *5*) *2 True*, *set-bit 5 0 False*, *set-bit* (− *5*) *1 False*
, *set-bit 5 32 True*, *set-bit 5 32 False*, *set-bit* (− *5*) *32 True*, *set-bit* (− *5*) *32 False*
, *1* << *2*, −*1* << *3*, *1* << *32*, *1* << *0*
, *100* >> *3*, −*100* >> *3*, *100* >> *32*, −*100* >> *32*
, *100* >>> *3*, −*100* >>> *3*, *100* >>> *32*, −*100* >>> *32*] :: *uint32 list*)
=
[*1*, *4294967295*, *5*, *4294967295*, *305419896*
, *18*
, *126*
, *108*
, *4294967205*
, *11*, *1*, *4294967295*, *4294967285*, *0*
, *2*, *4294967294*
, *15*, *4294967281*, *20*, *1891143032*
, *1*, *1431655763*, *0*, *0*
, *2*, *2*, *4294967291*, *5*
, *21*, *4294967295*, *4*, *4294967289*
, *5*, *5*, *4294967291*, *4294967291*
, *4*, *4294967288*, *0*, *1*
, *12*, *536870899*, *0*, *0*
, *12*, *4294967283*, *0*, *4294967295*]) ∧
([(*0x5* :: *uint32*) = *0x5*, (*0x5* :: *uint32*) = *0x6*
, (*0x5* :: *uint32*) < *0x5*, (*0x5* :: *uint32*) < *0x6*, (−*5* :: *uint32*) < *6*, (*6* :: *uint32*) < −*5*
, (*0x5* :: *uint32*) ≤ *0x5*, (*0x5* :: *uint32*) ≤ *0x4*, (−*5* :: *uint32*) ≤ *6*, (*6* :: *uint32*) ≤ −*5*
, (*0x7FFFFFFF* :: *uint32*) < *0x80000000*, (*0xFFFFFFFF* :: *uint32*) < *0*, (*0x80000000* :: *uint32*) < *0x7FFFFFFF*
, *bit* (*0x7FFFFFFF* :: *uint32*) *0*, *bit* (*0x7FFFFFFF* :: *uint32*) *31*, *bit* (*0x80000000* :: *uint32*) *31*, *bit* (*0x80000000* :: *uint32*) *32*
]
=
[*True*, *False*
, *False*, *True*, *False*, *True*
, *True*, *False*, *False*, *True*
, *True*, *False*, *False*
, *True*, *False*, *True*, *False*
]) ∧
([*integer-of-uint32 0*, *integer-of-uint32 0x7FFFFFFF*, *integer-of-uint32 0x80000000*, *integer-of-uint32 0xAAAAAAAA*]
=
[*0*, *0x7FFFFFFF*, *0x80000000*, *0xAAAAAAAA*])

export-code *test-uint32* **checking** *SML Haskell? OCaml? Scala*

notepad begin
⟨*proof*⟩
end

definition *test-uint32'* :: *uint32*
where *test-uint32' = 0 + 10 − 14 * 3 div 6 mod 3 << 3 >> 2*
⟨*ML*⟩

lemma *x AND y = x OR* (*y* :: *uint32*)
quickcheck[*random*, *expect=counterexample*]
quickcheck[*exhaustive*, *expect=counterexample*]
⟨*proof*⟩

lemma (*x* :: *uint32*) *AND x = x OR x*
quickcheck[*narrowing*, *expect=no-counterexample*]
⟨*proof*⟩

lemma (*f* :: *uint32* ⇒ *unit*) = *g*
quickcheck[*narrowing*, *size=3*, *expect=no-counterexample*]
⟨*proof*⟩

end

12.2 Tests for *uint16*

context
includes *bit-operations-syntax*
begin

abbreviation (*input*) *sshiftr-uint16* (**infixl** *>>>* *55*)
where ‹*w >>> n* ≡ *signed-drop-bit-uint16 n w*›

definition *test-uint16* **where**
test-uint16 ⟷
(([*0x10001*, *−1*, *−65535*, *0xFFFF*, *0x1234*
, *0x5A AND 0x36*
, *0x5A OR 0x36*
, *0x5A XOR 0x36*
, *NOT 0x5A*
, *5 + 6*, *−5 + 6*, *−6 + 5*, *−5 + −6*, *0xFFFF + 1*
, *5 − 3*, *3 − 5*
, *5 * 3*, *−5 * 3*, *−5 * −4*, *0x1234 * 0x8765*
, *5 div 3*, *−5 div 3*, *−5 div −3*, *5 div −3*
, *5 mod 3*, *−5 mod 3*, *−5 mod −3*, *5 mod −3*
, *set-bit 5 4 True*, *set-bit* (*− 5*) *2 True*, *set-bit 5 0 False*, *set-bit* (*− 5*) *1 False*
, *set-bit 5 32 True*, *set-bit 5 32 False*, *set-bit* (*− 5*) *32 True*, *set-bit* (*− 5*) *32 False*
, *1 << 2*, *−1 << 3*, *1 << 16*, *1 << 0*
, *100 >> 3*, *−100 >> 3*, *100 >> 16*, *−100 >> 16*
, *100 >>> 3*, *−100 >>> 3*, *100 >>> 16*, *−100 >>> 16*] :: *uint16 list*)
=
[*1*, *65535*, *1*, *65535*, *4660*

```
   , 18
   , 126
   , 108
   , 65445
   , 11, 1, 65535, 65525, 0
   , 2, 65534
   , 15, 65521, 20, 39556
   , 1, 21843, 0, 0
   , 2, 2, 65531, 5
   , 21, 65535, 4, 65529
   , 5, 5, 65531, 65531
   , 4, 65528, 0, 1
   , 12, 8179, 0, 0
   , 12, 65523, 0, 65535]) ∧
  ([ (0x5 :: uint16) = 0x5, (0x5 :: uint16) = 0x6
  , (0x5 :: uint16) < 0x5, (0x5 :: uint16) < 0x6, (−5 :: uint16) < 6, (6 :: uint16)
< −5
  , (0x5 :: uint16) ≤ 0x5, (0x5 :: uint16) ≤ 0x4, (−5 :: uint16) ≤ 6, (6 :: uint16)
≤ −5
  , (0x7FFF :: uint16) < 0x8000, (0xFFFF :: uint16) < 0, (0x8000 :: uint16) <
0x7FFF
   , bit (0x7FFF :: uint16) 0, bit (0x7FFF :: uint16) 15, bit (0x8000 :: uint16)
15, bit (0x8000 :: uint16) 16
  ]
  =
  [ True, False
  , False, True, False, True
  , True, False, False, True
  , True, False, False
  , True, False, True, False
  ]) ∧
  ([integer-of-uint16 0, integer-of-uint16 0x7FFF, integer-of-uint16 0x8000, inte-
ger-of-uint16 0xAAAA]
  =
  [0, 0x7FFF, 0x8000, 0xAAAA])

export-code test-uint16 checking Haskell? Scala
export-code test-uint16 checking SML-word

notepad begin
⟨proof⟩
end

lemma (x :: uint16) AND x = x OR x
quickcheck[narrowing, expect=no-counterexample]
⟨proof⟩

lemma (f :: uint16 ⇒ unit) = g
quickcheck[narrowing, size=3, expect=no-counterexample]
```

⟨*proof*⟩

end

12.3 Tests for *uint8*

context
includes *bit-operations-syntax*
begin

abbreviation (*input*) *sshiftr-uint8* (**infixl** >>> 55)
where ‹*w* >>> *n* ≡ *signed-drop-bit-uint8 n w*›

definition *test-uint8* **where**
test-uint8 ⟷
(([*0x101*, −*1*, −*255*, *0xFF*, *0x12*
, *0x5A AND 0x36*
, *0x5A OR 0x36*
, *0x5A XOR 0x36*
, *NOT 0x5A*
, *5* + *6*, −*5* + *6*, −*6* + *5*, −*5* + −*6*, *0xFF* + *1*
, *5* − *3*, *3* − *5*
, *5* ∗ *3*, −*5* ∗ *3*, −*5* ∗ −*4*, *0x12* ∗ *0x87*
, *5 div 3*, −*5 div 3*, −*5 div* −*3*, *5 div* −*3*
, *5 mod 3*, −*5 mod 3*, −*5 mod* −*3*, *5 mod* −*3*
, *set-bit 5 4 True*, *set-bit* (− *5*) *2 True*, *set-bit 5 0 False*, *set-bit* (− *5*) *1 False*
, *set-bit 5 32 True*, *set-bit 5 32 False*, *set-bit* (− *5*) *32 True*, *set-bit* (− *5*) *32 False*
, *1* << *2*, −*1* << *3*, *1* << *8*, *1* << *0*
, *100* >> *3*, −*100* >> *3*, *100* >> *8*, −*100* >> *8*
, *100* >>> *3*, −*100* >>> *3*, *100* >>> *8*, −*100* >>> *8*] :: *uint8 list*)
=
[*1*, *255*, *1*, *255*, *18*
, *18*
, *126*
, *108*
, *165*
, *11*, *1*, *255*, *245*, *0*
, *2*, *254*
, *15*, *241*, *20*, *126*
, *1*, *83*, *0*, *0*
, *2*, *2*, *251*, *5*
, *21*, *255*, *4*, *249*
, *5*, *5*, *251*, *251*
, *4*, *248*, *0*, *1*
, *12*, *19*, *0*, *0*
, *12*, *243*, *0*, *255*]) ∧
([(*0x5* :: *uint8*) = *0x5*, (*0x5* :: *uint8*) = *0x6*
, (*0x5* :: *uint8*) < *0x5*, (*0x5* :: *uint8*) < *0x6*, (−*5* :: *uint8*) < *6*, (*6* :: *uint8*) <

-5
, $(0x5 :: uint8) \leq 0x5$, $(0x5 :: uint8) \leq 0x4$, $(-5 :: uint8) \leq 6$, $(6 :: uint8) \leq -5$
, $(0x7F :: uint8) < 0x80$, $(0xFF :: uint8) < 0$, $(0x80 :: uint8) < 0x7F$
, $bit\ (0x7F :: uint8)\ 0$, $bit\ (0x7F :: uint8)\ 7$, $bit\ (0x80 :: uint8)\ 7$, $bit\ (0x80 :: uint8)\ 8$
]
=
[*True*, *False*
, *False*, *True*, *False*, *True*
, *True*, *False*, *False*, *True*
, *True*, *False*, *False*
, *True*, *False*, *True*, *False*
]) ∧
([*integer-of-uint8 0*, *integer-of-uint8 0x7F*, *integer-of-uint8 0x80*, *integer-of-uint8 0xAA*]
=
[*0*, *0x7F*, *0x80*, *0xAA*])

export-code *test-uint8* **checking** *SML Haskell? Scala*

notepad begin

⟨*proof*⟩
end
⟨*ML*⟩

definition *test-uint8′* :: *uint8*
where *test-uint8′ = 0 + 10 − 14 ∗ 3 div 6 mod 3 << 3 >> 2*
⟨*ML*⟩

lemma *x AND y = x OR* (*y* :: *uint8*)
quickcheck[*random*, *expect=counterexample*]
quickcheck[*exhaustive*, *expect=counterexample*]
⟨*proof*⟩

lemma (*x* :: *uint8*) *AND x = x OR x*
quickcheck[*narrowing*, *expect=no-counterexample*]
⟨*proof*⟩

lemma (*f* :: *uint8* ⇒ *unit*) = *g*
quickcheck[*narrowing*, *size=3*, *expect=no-counterexample*]
⟨*proof*⟩

12.4 Tests for *uint*

context
includes *bit-operations-syntax*
begin

abbreviation (*input*) *sshiftr-uint* (**infixl** *>>> 55*)
where ‹*w >>> n ≡ signed-drop-bit-uint n w*›

definition *test-uint ≡ let*
test-list1 = (let
HS = uint-of-int (2 ^ (dflt-size − 1))
in
([HS + HS + 1, −1, −HS − HS + 5, HS + (HS − 1), 0x12
, 0x5A AND 0x36
, 0x5A OR 0x36
, 0x5A XOR 0x36
, NOT 0x5A
, 5 + 6, −5 + 6, −6 + 5, −5 + −6, HS + (HS − 1) + 1
, 5 − 3, 3 − 5
*, 5 * 3, −5 * 3, −5 * −4, 0x12345678 * 0x87654321*]
@ (if dflt-size > 4 then
[5 div 3, −5 div 3, −5 div −3, 5 div −3
, 5 mod 3, −5 mod 3, −5 mod −3, 5 mod −3
, set-bit 5 4 True, set-bit (− 5) 2 True, set-bit 5 0 False, set-bit (− 5) 1 False
, set-bit 5 dflt-size True, set-bit 5 dflt-size False, set-bit (− 5) dflt-size True,
set-bit (− 5) dflt-size False
, 1 << 2, −1 << 3, push-bit dflt-size 1, 1 << 0
, 31 >> 3, −1 >> 3, 31 >> dflt-size, −1 >> dflt-size
, 15 >>> 2, −1 >>> 3, 15 >>> dflt-size, −1 >>> dflt-size]
else []) :: uint list));

test-list2 = (let
S = wivs-shift
in
([1, −1, −S + 5, S − 1, 0x12
, 0x5A AND 0x36
, 0x5A OR 0x36
, 0x5A XOR 0x36
, NOT 0x5A
, 5 + 6, −5 + 6, −6 + 5, −5 + −6, 0
, 5 − 3, 3 − 5
*, 5 * 3, −5 * 3, −5 * −4, 0x12345678 * 0x87654321*]
@ (if dflt-size > 4 then
[5 div 3, (S − 5) div 3, (S − 5) div (S − 3), 5 div (S − 3)
, 5 mod 3, (S − 5) mod 3, (S − 5) mod (S − 3), 5 mod (S − 3)
, set-bit 5 4 True, −1, set-bit 5 0 False, −7
, 5, 5, −5, −5
, 4, −8, 0, 1
, 3, (S >> 3) − 1, 0, 0
, 3, (S >> 1) + (S >> 1) − 1, 0, −1]
else []) :: int list));

test-list-c1 = (let

```
    HS = uint-of-int ((2^(dflt-size − 1)))
  in
 [ (0x5 :: uint) = 0x5, (0x5 :: uint) = 0x6
  , (0x5 :: uint) < 0x5, (0x5 :: uint) < 0x6, (−5 :: uint) < 6, (6 :: uint) < −5
  , (0x5 :: uint) ≤ 0x5, (0x5 :: uint) ≤ 0x4, (−5 :: uint) ≤ 6, (6 :: uint) ≤ −5
  , (HS − 1) < HS, (HS + HS − 1) < 0, HS < HS − 1
  , bit (HS − 1) 0, bit (HS − 1 :: uint) (dflt-size − 1), bit (HS :: uint) (dflt-size
− 1), bit (HS :: uint) dflt-size
  ]);

 test-list-c2 =
  [ True, False
  , False, dflt-size≥2, dflt-size=3, dflt-size≠3
  , True, False, dflt-size=3, dflt-size≠3
  , True, False, False
  , dflt-size≠1, False, True, False
  ]
in
 test-list1 = map uint-of-int test-list2
∧ test-list-c1 = test-list-c2
```

export-code *test-uint* **checking** *SML Haskell? OCaml? Scala*

lemma *test-uint*
quickcheck[*exhaustive*, *expect=no-counterexample*]
⟨*proof*⟩

lemma *x AND y = x OR* (*y* :: *uint*)
quickcheck[*random*, *expect=counterexample*]
quickcheck[*exhaustive*, *expect=counterexample*]
⟨*proof*⟩

lemma (*x* :: *uint*) *AND x = x OR x*
quickcheck[*narrowing*, *expect=no-counterexample*]
⟨*proof*⟩

lemma (*f* :: *uint* ⇒ *unit*) = *g*
quickcheck[*narrowing*, *size=3*, *expect=no-counterexample*]
⟨*proof*⟩

12.5 Tests for *uint64*

context
 includes *bit-operations-syntax*
begin

abbreviation (*input*) *sshiftr-uint64* (**infixl** >>> *55*)
 where ‹*w* >>> *n* ≡ *signed-drop-bit-uint64 n w*›

definition *test-uint64* **where**
test-uint64 ⟷
(([*0x10000000000000001*, *−1*, *−9223372036854775808*, *0xFFFFFFFFFFFFFFFF*, *0x1234567890ABCDEF*
, *0x5A AND 0x36*
, *0x5A OR 0x36*
, *0x5A XOR 0x36*
, *NOT 0x5A*
, *5 + 6*, *−5 + 6*, *−6 + 5*, *−5 + (− 6)*, *0xFFFFFFFFFFFFFFFFFFFF + 1*
, *5 − 3*, *3 − 5*
, *5 ∗ 3*, *−5 ∗ 3*, *−5 ∗ −4*, *0x1234567890ABCDEF ∗ 0xFEDCBA0987654321*
, *5 div 3*, *−5 div 3*, *−5 div −3*, *5 div −3*
, *5 mod 3*, *−5 mod 3*, *−5 mod −3*, *5 mod −3*
, *set-bit 5 4 True*, *set-bit (− 5) 2 True*, *set-bit 5 0 False*, *set-bit (− 5) 1 False*
, *set-bit 5 64 True*, *set-bit 5 64 False*, *set-bit (− 5) 64 True*, *set-bit (− 5) 64 False*
, *1 << 2*, *−1 << 3*, *1 << 64*, *1 << 0*
, *100 >> 3*, *−100 >> 3*, *100 >> 64*, *−100 >> 64*
, *100 >>> 3*, *−100 >>> 3*, *100 >>> 64*, *−100 >>> 64*] :: *uint64 list*)
=
[*1*, *18446744073709551615*, *9223372036854775808*, *18446744073709551615*, *1311768467294899695*
, *18*
, *126*
, *108*
, *18446744073709551525*
, *11*, *1*, *18446744073709551615*, *18446744073709551605*, *0*
, *2*, *18446744073709551614*
, *15*, *18446744073709551601*, *20*, *14000077364136384719*
, *1*, *6148914691236517203*, *0*, *0*
, *2*, *2*, *18446744073709551611*, *5*
, *21*, *18446744073709551615*, *4*, *18446744073709551609*
, *5*, *5*, *18446744073709551611*, *18446744073709551611*
, *4*, *18446744073709551608*, *0*, *1*
, *12*, *2305843009213693939*, *0*, *0*
, *12*, *18446744073709551603*, *0*, *18446744073709551615*]) ∧
([(*0x5* :: *uint64*) = *0x5*, (*0x5* :: *uint64*) = *0x6*
, (*0x5* :: *uint64*) < *0x5*, (*0x5* :: *uint64*) < *0x6*, (*−5* :: *uint64*) < *6*, (*6* :: *uint64*) < *−5*
, (*0x5* :: *uint64*) ≤ *0x5*, (*0x5* :: *uint64*) ≤ *0x4*, (*−5* :: *uint64*) ≤ *6*, (*6* :: *uint64*) ≤ *−5*
, (*0x7FFFFFFFFFFFFFFF* :: *uint64*) < *0x8000000000000000*, (*0xFFFFFFFFFFFFFFFF* :: *uint64*) < *0*, (*0x8000000000000000* :: *uint64*) < *0x7FFFFFFFFFFFFFFF*
, *bit* (*0x7FFFFFFFFFFFFFFF* :: *uint64*) *0*, *bit* (*0x7FFFFFFFFFFFFFFF* :: *uint64*) *63*, *bit* (*0x8000000000000000* :: *uint64*) *63*, *bit* (*0x8000000000000000* :: *uint64*) *64*
]
=
[*True*, *False*

, *False*, *True*, *False*, *True*
, *True*, *False*, *False*, *True*
, *True*, *False*, *False*
, *True*, *False*, *True*, *False*
]) ∧
([*integer-of-uint64 0*, *integer-of-uint64 0x7FFFFFFFFFFFFFFF*, *integer-of-uint64 0x8000000000000000*, *integer-of-uint64 0xAAAAAAAAAAAAAAAA*]
=
[*0*, *0x7FFFFFFFFFFFFFFF*, *0x8000000000000000*, *0xAAAAAAAAAAAAAAAA*])

value [*nbe*] [*0x10000000000000001*, *−1*, *−9223372036854775808*, *0xFFFFFFFFFFFFFFFF*, *0x1234567890ABCDEF*
, *0x5A AND 0x36*
, *0x5A OR 0x36*
, *0x5A XOR 0x36*
, *NOT 0x5A*
, *5 + 6*, *−5 + 6*, *−6 + 5*, *−5 + (− 6)*, *0xFFFFFFFFFFFFFFFFFF + 1*
, *5 − 3*, *3 − 5*
, *5 ∗ 3*, *−5 ∗ 3*, *−5 ∗ −4*, *0x1234567890ABCDEF ∗ 0xFEDCBA0987654321*
, *5 div 3*, *−5 div 3*, *−5 div −3*, *5 div −3*
, *5 mod 3*, *−5 mod 3*, *−5 mod −3*, *5 mod −3*
, *set-bit 5 4 True*, *set-bit (− 5) 2 True*, *set-bit 5 0 False*, *set-bit (− 5) 1 False*
, *set-bit 5 64 True*, *set-bit 5 64 False*, *set-bit (− 5) 64 True*, *set-bit (− 5) 64 False*
, *1 << 2*, *−1 << 3*, *1 << 64*, *1 << 0*
, *100 >> 3*, *−100 >> 3*, *100 >> 64*, *−100 >> 64*
, *100 >>> 3*, *−100 >>> 3*, *100 >>> 64*, *−100 >>> 64*] :: *uint64 list*

export-code *test-uint64* **checking** *SML Haskell? OCaml? Scala*

notepad begin
⟨*proof*⟩
end
⟨*ML*⟩

definition *test-uint64′* :: *uint64*
where *test-uint64′ = 0 + 10 − 14 ∗ 3 div 6 mod 3 << 3 >> 2*

end

12.6 Tests for casts

definition *test-casts* :: *bool*
where *test-casts* ⟷
map uint8-of-uint32 [*10*, *0*, *0xFE*, *0xFFFFFFFF*] = [*10*, *0*, *0xFE*, *0xFF*] ∧
map uint8-of-uint64 [*10*, *0*, *0xFE*, *0xFFFFFFFFFFFFFFFF*] = [*10*, *0*, *0xFE*, *0xFF*] ∧
map uint32-of-uint8 [*10*, *0*, *0xFF*] = [*10*, *0*, *0xFF*] ∧
map uint64-of-uint8 [*10*, *0*, *0xFF*] = [*10*, *0*, *0xFF*]

definition *test-casts'* :: *bool*
where *test-casts'* ⟷
map uint8-of-uint16 [*10*, *0*, *0xFE*, *0xFFFF*] = [*10*, *0*, *0xFE*, *0xFF*] ∧
map uint16-of-uint8 [*10*, *0*, *0xFF*] = [*10*, *0*, *0xFF*] ∧
map uint16-of-uint32 [*10*, *0*, *0xFFFE*, *0xFFFFFFFF*] = [*10*, *0*, *0xFFFE*, *0xFFFF*]
∧
map uint16-of-uint64 [*10*, *0*, *0xFFFE*, *0xFFFFFFFFFFFFFFFF*] = [*10*, *0*,
0xFFFE, *0xFFFF*] ∧
map uint32-of-uint16 [*10*, *0*, *0xFFFF*] = [*10*, *0*, *0xFFFF*] ∧
map uint64-of-uint16 [*10*, *0*, *0xFFFF*] = [*10*, *0*, *0xFFFF*]

definition *test-casts''* :: *bool*
where *test-casts''* ⟷
map uint32-of-uint64 [*10*, *0*, *0xFFFFFFFE*, *0xFFFFFFFFFFFFFFFF*] = [*10*,
0, *0xFFFFFFFE*, *0xFFFFFFFF*] ∧
map uint64-of-uint32 [*10*, *0*, *0xFFFFFFFF*] = [*10*, *0*, *0xFFFFFFFF*]

export-code *test-casts test-casts''* **checking** *SML Haskell? Scala*
export-code *test-casts''* **checking** *OCaml?*
export-code *test-casts'* **checking** *Haskell? Scala*

notepad begin
⟨*proof*⟩
end
⟨*ML*⟩

definition *test-casts-uint* :: *bool* **where**
test-casts-uint ⟷
map uint-of-uint32 ([*0*, *10*] @ (*if dflt-size* < *32 then* [*push-bit* (*dflt-size* − *1*) *1*,
0xFFFFFFFF] *else* [*0xFFFFFFFF*])) =
[*0*, *10*] @ (*if dflt-size* < *32 then* [*push-bit* (*dflt-size* − *1*) *1*, (*push-bit dflt-size 1*)
− *1*] *else* [*0xFFFFFFFF*]) ∧
map uint32-of-uint [*0*, *10*, *if dflt-size* < *32 then push-bit* (*dflt-size* − *1*) *1 else*
0xFFFFFFFF] =
[*0*, *10*, *if dflt-size* < *32 then push-bit* (*dflt-size* − *1*) *1 else 0xFFFFFFFF*] ∧
map uint-of-uint64 [*0*, *10*, *push-bit* (*dflt-size* − *1*) *1*, *0xFFFFFFFFFFFFFFFF*]
=
[*0*, *10*, *push-bit* (*dflt-size* − *1*) *1*, (*push-bit dflt-size 1*) − *1*] ∧
map uint64-of-uint [*0*, *10*, *push-bit* (*dflt-size* − *1*) *1*] =
[*0*, *10*, *push-bit* (*dflt-size* − *1*) *1*]

definition *test-casts-uint'* :: *bool* **where**
test-casts-uint' ⟷
map uint-of-uint16 [*0*, *10*, *0xFFFF*] = [*0*, *10*, *0xFFFF*] ∧
map uint16-of-uint [*0*, *10*, *0xFFFF*] = [*0*, *10*, *0xFFFF*]

definition *test-casts-uint''* :: *bool* **where**
test-casts-uint'' ⟷

$map\ uint\text{-}of\text{-}uint8\ [0,\ 10,\ 0xFF] = [0,\ 10,\ 0xFF] \wedge$
$map\ uint8\text{-}of\text{-}uint\ [0,\ 10,\ 0xFF] = [0,\ 10,\ 0xFF]$

end

end

end

Chapter 13

Implementation of bit operations on int by target language operations

theory *Code-Target-Bits-Int*
imports
Bits-Integer
HOL−Library.Code-Target-Int
begin

context
includes *bit-operations-syntax*
begin

declare [[*code drop*:
(*AND*) :: *int* ⇒ - (*OR*) :: *int* ⇒ - (*XOR*) :: *int* ⇒ - *NOT* :: *int* ⇒ -
lsb :: *int* ⇒ - *set-bit* :: *int* ⇒ - *bit* :: *int* ⇒ -
push-bit :: - ⇒ *int* ⇒ - *drop-bit* :: - ⇒ *int* ⇒ -
int-of-integer-symbolic
]]

lemma [*code-unfold*]:
‹*of-bool* (*odd i*) = *i AND 1*› **for** *i* :: *int*
⟨*proof*⟩

lemma [*code-unfold*]:
‹*bit x n* ⟷ *x AND* (*push-bit n 1*) ≠ *0*› **for** *x* :: *int*
⟨*proof*⟩

context
includes *integer.lifting*
begin

lemma *bit-int-code* [*code*]:
bit (*int-of-integer x*) *n* = *bit x n*
⟨*proof*⟩

lemma *and-int-code* [*code*]:
int-of-integer i AND int-of-integer j = *int-of-integer* (*i AND j*)
⟨*proof*⟩

lemma *or-int-code* [*code*]:
int-of-integer i OR int-of-integer j = *int-of-integer* (*i OR j*)
⟨*proof*⟩

lemma *xor-int-code* [*code*]:
int-of-integer i XOR int-of-integer j = *int-of-integer* (*i XOR j*)
⟨*proof*⟩

lemma *not-int-code* [*code*]:
NOT (*int-of-integer i*) = *int-of-integer* (*NOT i*)
⟨*proof*⟩

lemma *push-bit-int-code* [*code*]:
‹*push-bit n* (*int-of-integer x*) = *int-of-integer* (*push-bit n x*)›
⟨*proof*⟩

lemma *drop-bit-int-code* [*code*]:
‹*drop-bit n* (*int-of-integer x*) = *int-of-integer* (*drop-bit n x*)›
⟨*proof*⟩

lemma *take-bit-int-code* [*code*]:
‹*take-bit n* (*int-of-integer x*) = *int-of-integer* (*take-bit n x*)›
⟨*proof*⟩

lemma *lsb-int-code* [*code*]:
lsb (*int-of-integer x*) = *lsb x*
⟨*proof*⟩

lemma *set-bit-int-code* [*code*]:
set-bit (*int-of-integer x*) *n b* = *int-of-integer* (*set-bit x n b*)
⟨*proof*⟩

lemma *int-of-integer-symbolic-code* [*code*]:
int-of-integer-symbolic = *int-of-integer*
⟨*proof*⟩

context
begin

qualified definition *even* :: ‹*int* ⇒ *bool*›
where [*code-abbrev*]: ‹*even* = *Parity.even*›

end

lemma [*code*]:
‹*Code-Target-Bits-Int.even i* ⟷ *i AND 1 = 0*›
⟨*proof*⟩

lemma *bin-rest-code*:
int-of-integer i div 2 = int-of-integer (*bin-rest-integer i*)
⟨*proof*⟩

end

end

end

theory *Native-Word-Test-Emu* **imports**
Native-Word-Test
Code-Target-Bits-Int
begin

13.1 Test cases for emulation of native words

13.1.1 Tests for *uint16*

Test that *uint16* is emulated for PolyML and OCaml via *16 word* if *Native-Word.Code-Target-Bits-Int* is imported.

definition *test-uint16-emulation* :: *bool* **where**
test-uint16-emulation ⟷ (*0xFFFFF − 0x1000* = (*0xEFFF* :: *uint16*))

export-code *test-uint16-emulation* **checking** *SML OCaml?*
— test the other target languages as well *Haskell? Scala*

notepad begin
⟨*proof*⟩
end

⟨*ML*⟩

lemma *x AND y = x OR* (*y* :: *uint16*)
quickcheck[*random*, *expect*=*counterexample*]
quickcheck[*exhaustive*, *expect*=*counterexample*]
⟨*proof*⟩

13.1.2 Tests for *uint8*

Test that *uint8* is emulated for OCaml via *8 word* if *Native-Word.Code-Target-Bits-Int* is imported.

definition *test-uint8-emulation* :: *bool* **where**
test-uint8-emulation $\longleftrightarrow$ (*0xFFF* − *0x10* = (*0xEF* :: *uint8*))

export-code *test-uint8-emulation* **checking** *OCaml?*
— test the other target languages as well *SML Haskell? Scala*

end

theory *Native-Word-Test-PolyML* **imports**
Native-Word-Test
begin

13.2 Test with PolyML

test-code
test-uint64 test-uint64′ = *0x12*
test-uint32 test-uint32′ = *0x12*
test-uint8 test-uint8′ = *0x12*
test-uint
test-casts test-casts″
test-casts-uint test-casts-uint″
in *PolyML*

end

theory *Native-Word-Test-PolyML2* **imports**
Native-Word-Test-Emu
begin

test-code
test-uint16 test-uint16-emulation
test-casts′
test-casts-uint′
in *PolyML*

end

theory *Native-Word-Test-PolyML64* **imports**
Native-Word-Test
begin

test-code *test-uint64′ = 0x12*
in *PolyML*

⟨*ML*⟩

end

theory *Native-Word-Test-Scala* **imports**
Native-Word-Test
begin

13.3 Test with Scala

In Scala, *uint* and *uint32* are both implemented as type `Int`. When they are used in the same generated program, we have to suppress the type class instances for one of them.

code-printing class-instance *uint32* :: *equal* ⇀ (*Scala*) –

test-code
test-uint64 test-uint64′ = 0x12
test-uint32 test-uint32′ = 0x12
test-uint16
test-uint8 test-uint8′ = 0x12
test-uint
test-casts test-casts′ test-casts″
test-casts-uint test-casts-uint′ test-casts-uint″
in *Scala*

end

Chapter 14

User guide for native words

This tutorial explains how to best use the types for native words like *uint32* in your formalisation. You can base your formalisation

1. either directly on these types,

2. or on the generic *'a word* and only introduce native words a posteriori via code generator refinement.

The first option causes the least overhead if you have to prove only little about the words you use and start a fresh formalisation. Just use the native type *uint32* instead of *32 word* and similarly for *uint64*, *uint16*, and *uint8*. As native word types are meant only for code generation, the lemmas about *'a word* have not been duplicated, but you can transfer theorems between native word types and *'a word* using the transfer package.

Note, however, that this option restricts your work a bit: your own functions cannot be "polymorphic" in the word length, but you have to define a separate function for every word length you need.

The second option is recommended if you already have a formalisation based on *'a word* or if your proofs involve words and their properties. It separates code generation from modelling and proving, i.e., you can work with words as usual. Consequently, you have to manually setup the code generator to use the native types wherever you want. The following describes how to achieve this with moderate effort.

Note, however, that some target languages of the code generator (especially OCaml) do not support all the native word types provided. Therefore, you should only import those types that you need – the theory file for each type mentions at the top the restrictions for code generation. For example, PolyML does not provide the Word16 structure, and OCaml provides neither Word8 nor Word16. You can still use these theories provided that you also import the theory *Native-Word.Code-Target-Bits-Int* (which implements *int*

by target-language integers), but these words will be implemented via Isabelle's *Word* library, i.e., you do not gain anything in terms of efficiency.

There is a separate code target *SML-word* **for SML.** If you use one of the native words that PolyML does not support (such as *uint16* and *uint64* in 32-bit mode), but would like to map its operations to the Standard Basis Library functions, make sure to use the target *SML-word* instead of *SML*; if you only use native word sizes that PolyML supports, you can stick with *SML*. This ensures that code generation within Isabelle as used by *Quickcheck*, *value* and @{code} in ML blocks continues to work.

14.1 Lifting functions from *'a word* to native words

This section shows how to convert functions from *'a word* to native words. For example, the following function *sum-squares* computes the sum of the first n square numbers in 16 bit arithmetic using a tail-recursive function *gen-sum-squares* with accumulator; for convenience, *sum-squares-int* takes an integer instead of a word.

```
function gen-sum-squares :: 16 word ⇒ 16 word ⇒ 16 word where
  gen-sum-squares accum n =
  (if n = 0 then accum else gen-sum-squares (accum + n * n) (n − 1))⟨proof⟩⟨proof⟩
definition sum-squares :: 16 word ⇒ 16 word where
   sum-squares = gen-sum-squares 0

definition sum-squares-int :: int ⇒ 16 word where
  sum-squares-int n = sum-squares (word-of-int n)
```

The generated code for *sum-squares* and *sum-squares-int* emulates words with unbounded integers and explicit modulus as specified in the theory *HOL−Library.Word*. But for efficiency, we want that the generated code uses machine words and machine arithmetic. Unfortunately, as *'a word* is polymorphic in the word length, the code generator can only do this if we use another type for machine words. The theory *Native-Word.Uint16* defines the type *uint16* for machine words of 16 bits. We just have to follow two steps to use it:

First, we lift all our functions from *16 word* to *uint16*, i.e., *sum-squares*, *gen-sum-squares*, and *sum-squares-int* in our case. The theory *Native-Word.Uint16* sets up the lifting package for this and has already taken care of the arithmetic and bit-wise operations.

```
lift-definition gen-sum-squares-uint :: uint16 ⇒ uint16 ⇒ uint16
  is gen-sum-squares ⟨proof⟩
lift-definition sum-squares-uint :: uint16 ⇒ uint16 is sum-squares ⟨proof⟩
lift-definition sum-squares-int-uint :: int ⇒ uint16 is sum-squares-int ⟨proof⟩
```

Second, we also have to transfer the code equations for our functions. The

attribute *Transfer.transferred* takes care of that, but it is better to check that the transfer succeeded: inspect the theorem to check that the new constants are used throughout.

lemmas [*Transfer.transferred*, *code*] =
gen-sum-squares.simps
sum-squares-def
sum-squares-int-def

Finally, we export the code to standard ML. We use the target *SML-word* instead of *SML* to have the operations on *uint16* mapped to the Standard Basis Library. As PolyML does not provide a Word16 type, the mapping for *uint16* is only active in the refined target *SML-word*.

export-code *sum-squares-int-uint* **in** *SML-word*

Nevertheless, we can still evaluate terms with *uint16* within Isabelle, i.e., PolyML, but this will be translated to *16 word* and therefore less efficient.

value *sum-squares-int-uint 40*

14.2 Storing native words in datatypes

The above lifting is necessary for all functions whose type mentions the word type. Fortunately, we do not have to duplicate functions that merely operate on datatypes that contain words. Nevertheless, we have to tell the code generator that these functions should call the new ones, which operate on machine words. This section shows how to achieve this with data refinement.

14.2.1 Example: expressions and two semantics

As the running example, we consider a language of expressions (literal values, less-than comparisions and conditional) where values are either booleans or 32-bit words. The original specification uses the type *32 word*.

datatype *val* = *Bool bool* | *Word 32 word*
datatype *expr* = *Lit val* | *LT expr expr* | *IF expr expr expr*

abbreviation (*input*) *word* :: *32 word* ⇒ *expr* **where** *word i* ≡ *Lit* (*Word i*)
abbreviation (*input*) *bool* :: *bool* ⇒ *expr* **where** *bool i* ≡ *Lit* (*Bool i*)

— Denotational semantics of expressions, *None* denotes a type error
fun *eval* :: *expr* ⇒ *val option* **where**
eval (*Lit v*) = *Some v*
| *eval* (*LT* e_1 e_2) =
(*case* (*eval* e_1, *eval* e_2)
of (*Some* (*Word* i_1), *Some* (*Word* i_2)) ⇒ *Some* (*Bool* (i_1 < i_2))
| - ⇒ *None*)
| *eval* (*IF* e_1 e_2 e_3) =

(*case eval* e_1 *of Some* (*Bool b*) ⇒ *if b then eval* e_2 *else eval* e_3
| - ⇒ *None*)

— Small-step semantics of expressions, it gets stuck upon type errors.
inductive *step* :: *expr* ⇒ *expr* ⇒ *bool* (- → - [*50*, *50*] *60*) **where**
$e \to e' \Longrightarrow LT\ e\ e_2 \to LT\ e'\ e_2$
| $e \to e' \Longrightarrow LT$ (*word i*) $e \to LT$ (*word i*) e'
| LT (*word* i_1) (*word* i_2) → *bool* ($i_1 < i_2$)
| $e \to e' \Longrightarrow IF\ e\ e_1\ e_2 \to IF\ e'\ e_1\ e_2$
| IF (*bool True*) $e_1\ e_2 \to e_1$
| IF (*bool False*) $e_1\ e_2 \to e_2$

— Compile the inductive definition with the predicate compiler
code-pred (*modes*: $i \Rightarrow o \Rightarrow$ *bool as reduce*, $i \Rightarrow i \Rightarrow$ *bool as step'*) *step* ⟨*proof*⟩

14.2.2 Change the datatype to use machine words

Now, we want to use *uint32* instead of *32 word*. The goal is to make the code generator use the new type without duplicating any of the types (*val*, *expr*) or the functions (*eval*, *reduce*) on such types.

The constructor *Word* has *32 word* in its type, so we have to lift it to *Word'*, and the same holds for the case combinator *case-val*, which *case-val'* replaces.[1] Next, we set up the code generator accordingly: *Bool* and *Word'* are the new constructors for *val*, and *case-val'* is the new case combinator with an appropriate case certificate.[2] We delete the code equations for the old constructor *Word* and case combinator *case-val* such that the code generator reports missing adaptations.

lift-definition *Word'* :: *uint32* ⇒ *val* **is** *Word* ⟨*proof*⟩

code-datatype *Bool Word'*

[1] Note that we should not declare a case translation for the new case combinator because this will break parsing case expressions with old case combinator.

[2] Case certificates tell the code generator to replace the HOL case combinator for a datatype with the case combinator of the target language. Without a case certificate, the code generator generates a function that re-implements the case combinator; in a strict languages like ML or Scala, this means that the code evaluates all possible cases before it decides which one is taken.

Case certificates are described in Haftmann's PhD thesis [1, Def. 27]. For a datatype *dt* with constructors C_1 to C_n where each constructor C_i takes k_i parameters, the certificate for the case combinator *case-dt* looks as follows:

lemma
assumes CASE ≡ dt-case $c_1\ c_2\ \ldots\ c_n$
shows (CASE ($C_1\ a_{11}\ a_{12}\ \ldots\ a_{1k_1}$) ≡ $c_1\ a_{11}\ a_{12}\ \ldots\ a_{1k_1}$)
&&& (CASE ($C_2\ a_{21}\ a_{22}\ \ldots\ a_{2k_2}$) ≡ $c_2\ a_{21}\ a_{22}\ \ldots\ a_{2k_2}$)
&&& ...
&&& (CASE ($C_n\ a_{n1}\ a_{n2}\ \ldots\ a_{nk_n}$) ≡ $c_n\ a_{n1}\ a_{n2}\ \ldots\ a_{nk_n}$)

lift-definition *case-val′* :: (*bool* ⇒ *′a*) ⇒ (*uint32* ⇒ *′a*) ⇒ *val* ⇒ *′a* **is** *case-val* ⟨*proof*⟩

lemmas [*code*, *simp*] = *val.case* [*Transfer.transferred*]

lemma *case-val′-cert*:
fixes *bool word′ b w*
assumes *CASE* ≡ *case-val′ bool word′*
shows (*CASE* (*Bool b*) ≡ *bool b*) &&& (*CASE* (*Word′ w*) ≡ *word′ w*)
⟨*proof*⟩

⟨*ML*⟩

declare [[*code drop*: *case-val Word*]]

14.2.3 Make functions use functions on machine words

Finally, we merely have to change the code equations to use the new functions that operate on *uint32*. As before, the attribute *Transfer.transferred* does the job. In our example, we adapt the equality test on *val* (code equations *val.eq.simps*) and the denotational and small-step semantics (code equations *eval.simps* and *step.equation*, respectively).

We check that the adaptation has suceeded by exporting the functions. As we only use native word sizes that PolyML supports, we can use the usual target *SML* instead of *SML-word*.

lemmas [*code*] =
val.eq.simps[*THEN meta-eq-to-obj-eq*, *Transfer.transferred*, *THEN eq-reflection*]
eval.simps[*Transfer.transferred*]
step.equation[*Transfer.transferred*]

export-code *reduce step′ eval* **checking** *SML*

14.3 Troubleshooting

This section explains some possible problems when using native words. If you experience other difficulties, please contact the author.

14.3.1 *export-code* raises an exception

Probably, you have defined and are using a function on a native word type, but the code equation refers to emulated words. For example, the following defines a function *double* that doubles a word. When we try to export code for *double* without any further setup, *export-code* will raise an exception or generate code that does not compile.

lift-definition *double* :: *uint32* ⇒ *uint32* **is** λ*x*. *x* + *x* ⟨*proof*⟩

We have to prove a code equation that only uses the existing operations on *uint32*. Then, *export-code* works again.

lemma *double-code* [*code*]: *double* $n = n + n$
⟨*proof*⟩

14.3.2 The generated code does not compile

Probably, you have been exporting to a target language for which there is no setup, or your compiler does not provide the required API. Every theory for native words mentions at the start the limitations on code generation. Check that your concrete application meets all the requirements.

Alternatively, this might be an instance of the problem described in §14.3.1.

For Haskell, you have to enable the extension TypeSynonymInstances with `-XTypeSynonymInstances` if you are using polymorphic bit operations on the native word types.

14.3.3 The generated code is too slow

The generated code will most likely not be as fast as a direct implementation in the target language with manual tuning. This is because we want the configuration of the code generation to be sound (as it can be used to prove theorems in Isabelle). Therefore, the bit operations sometimes perform range checks before they call the target language API. Here are some examples:

- Shift distances and bit indices in target languages are often expected to fit into a bounded integer or word. However, the size of these types varies across target languages and platforms. Hence, no Isabelle/HOL type can model uniformly all of them. Instead, the bit operations use arbitrary-precision integers for such quantities and check at run-time that the values fit into a bounded integer or word, respectively – if not, they raise an exception.
- Division and modulo operations explicitly test whether the divisor is 0 and return the HOL value of division by 0 in that case. This is necessary because some languages leave the behaviour of division by 0 unspecified.

If you have better ideas how to eliminate such checks and speed up the generated code without sacrificing soundness, please contact the author!

Bibliography

[1] F. Haftmann. *Code Generation form Specifications in Higher-Order-Logic.* PhD thesis, Institut für Informatik, Technische Universität München, 2009.