

Interval Temporal Logic on Natural Numbers

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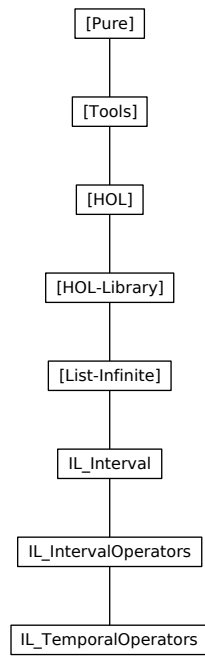
Abstract

We introduce a theory of temporal logic operators using sets of natural numbers as time domain, formalized in a shallow embedding manner. The theory comprises special natural intervals (theory `IL_Interval`: open and closed intervals, continuous and modulo intervals, interval traversing results), operators for shifting intervals to left/right on the number axis as well as expanding/contracting intervals by constant factors (theory `IL_IntervalOperators.thy`), and ultimately definitions and results for unary and binary temporal operators on arbitrary natural sets (theory `IL_TemporalOperators`).

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1 Intervals and operations for temporal logic declarations

```

theory IL-Interval
imports
  List-Infinite.InfiniteSet2
  List-Infinite.SetIntervalStep
begin

```

1.1 Time intervals – definitions and basic lemmata

1.1.1 Definitions

```

type-synonym Time = nat

```

```

type-synonym iT = Time set

```

Infinite interval starting at some natural n .

definition

```

iFROM :: Time  $\Rightarrow$  iT ([...]

```

where

```

[n...]  $\equiv$  {n.}

```

Finite interval starting at $0::'a$ and ending at some natural n .

definition

```

iTILL :: Time  $\Rightarrow$  iT ([...])

```

where

```

[...n]  $\equiv$  {..n}

```

Finite bounded interval containing the naturals between n and $n + d$. d denotes the difference between left and right interval bound. The number of elements is $d + (1::'a)$ so that an empty interval cannot be defined.

definition

```

iIN :: Time  $\Rightarrow$  nat  $\Rightarrow$  iT ([...])

```

where

```

[n...,d]  $\equiv$  {n..n+d}

```

Infinite modulo interval containing all naturals having the same division remainder modulo m as r , and beginning at n .

definition

```

iMOD :: Time  $\Rightarrow$  nat  $\Rightarrow$  iT ([ -, mod - ])

```

where

```

[r, mod m]  $\equiv$  { x. x mod m = r mod m  $\wedge$  r  $\leq$  x }

```

Finite bounded modulo interval containing all naturals having the same division remainder modulo m as r , beginning at n , and ending after c cycles

at $r + m * c$. The number of elements is $c + (1::'a)$ so that an empty interval cannot be defined.

definition

$iMODb :: Time \Rightarrow nat \Rightarrow nat \Rightarrow iT ([-, mod -, -])$

where

$[r, mod m, c] \equiv \{ x. x \text{ mod } m = r \text{ mod } m \wedge r \leq x \wedge x \leq r + m * c \}$

1.1.2 Membership in an interval

lemmas $iT-defs = iFROM-def iTILL-def iIN-def iMOD-def iMODb-def$

lemma $iFROM-iff: x \in [n..] = (n \leq x)$

$\langle proof \rangle$

lemma $iTILL-iff: x \in [..n] = (x \leq n)$

$\langle proof \rangle$

lemma $iIN-iff: x \in [n..,d] = (n \leq x \wedge x \leq n + d)$

$\langle proof \rangle$

lemma $iMOD-iff: x \in [r, mod m] = (x \text{ mod } m = r \text{ mod } m \wedge r \leq x)$

$\langle proof \rangle$

lemma $iMODb-iff: x \in [r, mod m, c] =$

$(x \text{ mod } m = r \text{ mod } m \wedge r \leq x \wedge x \leq r + m * c)$

$\langle proof \rangle$

lemma $iFROM-D: x \in [n..] \Longrightarrow (n \leq x)$

$\langle proof \rangle$

lemma $iTILL-D: x \in [..n] \Longrightarrow (x \leq n)$

$\langle proof \rangle$

corollary $iIN-geD: x \in [n..,d] \Longrightarrow n \leq x$

$\langle proof \rangle$

corollary $iIN-leD: x \in [n..,d] \Longrightarrow x \leq n + d$

$\langle proof \rangle$

corollary $iMOD-modD: x \in [r, mod m] \Longrightarrow x \text{ mod } m = r \text{ mod } m$

$\langle proof \rangle$

corollary $iMOD-geD: x \in [r, mod m] \Longrightarrow r \leq x$

$\langle proof \rangle$

corollary $iMODb-modD: x \in [r, mod m, c] \Longrightarrow x \text{ mod } m = r \text{ mod } m$

$\langle proof \rangle$

corollary $iMODb-geD: x \in [r, mod m, c] \Longrightarrow r \leq x$

$\langle proof \rangle$

corollary $iMODb-leD: x \in [r, mod m, c] \Longrightarrow x \leq r + m * c$

$\langle proof \rangle$

lemmas $iT-iff = iFROM-iff iTILL-iff iIN-iff iMOD-iff iMODb-iff$

lemmas $iT-drule =$

$iFROM-D$

$iTILL-D$

$iIN-geD iIN-leD$

$iMOD-modD iMOD-geD$

iMODb-modD iMODb-geD iMODb-leD

lemma

iFROM-I [intro]: $n \leq x \implies x \in [n..]$ **and**
iTILL-I [intro]: $x \leq n \implies x \in [..n]$ **and**
iIN-I [intro]: $n \leq x \implies x \leq n + d \implies x \in [n..,d]$ **and**
iMOD-I [intro]: $x \bmod m = r \bmod m \implies r \leq x \implies x \in [r, \bmod m]$ **and**
iMODb-I [intro]: $x \bmod m = r \bmod m \implies r \leq x \implies x \leq r + m * c \implies x \in [r, \bmod m, c]$
 ⟨proof⟩

lemma

iFROM-E [elim]: $x \in [n..] \implies (n \leq x \implies P) \implies P$ **and**
iTILL-E [elim]: $x \in [..n] \implies (x \leq n \implies P) \implies P$ **and**
iIN-E [elim]: $x \in [n..,d] \implies (n \leq x \implies x \leq n + d \implies P) \implies P$ **and**
iMOD-E [elim]: $x \in [r, \bmod m] \implies (x \bmod m = r \bmod m \implies r \leq x \implies P) \implies P$ **and**
iMODb-E [elim]: $x \in [r, \bmod m, c] \implies (x \bmod m = r \bmod m \implies r \leq x \implies x \leq r + m * c \implies P) \implies P$
 ⟨proof⟩

lemma *iIN-Suc-insert-conv*:

$\text{insert} (\text{Suc} (n + d)) [n..,d] = [n..,\text{Suc} d]$
 ⟨proof⟩

lemma *iTILL-Suc-insert-conv*: $\text{insert} (\text{Suc} n) [..n] = [..\text{Suc} n]$

⟨proof⟩

lemma *iMODb-Suc-insert-conv*:

$\text{insert} (r + m * \text{Suc} c) [r, \bmod m, c] = [r, \bmod m, \text{Suc} c]$
 ⟨proof⟩

lemma *iFROM-pred-insert-conv*: $\text{insert} (n - \text{Suc} 0) [n..] = [n - \text{Suc} 0..]$

⟨proof⟩

lemma *iIN-pred-insert-conv*:

$0 < n \implies \text{insert} (n - \text{Suc} 0) [n..,d] = [n - \text{Suc} 0..,\text{Suc} d]$
 ⟨proof⟩

lemma *iMOD-pred-insert-conv*:

$m \leq r \implies \text{insert} (r - m) [r, \bmod m] = [r - m, \bmod m]$
 ⟨proof⟩

lemma *iMODb-pred-insert-conv*:

$m \leq r \implies \text{insert} (r - m) [r, \bmod m, c] = [r - m, \bmod m, \text{Suc} c]$

<proof>

lemma *iFROM-Suc-pred-insert-conv*: $\text{insert } n \text{ [Suc } n \dots] = [n \dots]$

<proof>

lemma *iIN-Suc-pred-insert-conv*: $\text{insert } n \text{ [Suc } n \dots, d] = [n \dots, \text{Suc } d]$

<proof>

lemma *iMOD-Suc-pred-insert-conv*: $\text{insert } r \text{ [} r + m, \text{ mod } m] = [r, \text{ mod } m]$

<proof>

lemma *iMODb-Suc-pred-insert-conv*: $\text{insert } r \text{ [} r + m, \text{ mod } m, c] = [r, \text{ mod } m, \text{Suc } c]$

<proof>

lemmas *iT-Suc-insert* =

iIN-Suc-insert-conv

iTILL-Suc-insert-conv

iMODb-Suc-insert-conv

lemmas *iT-pred-insert* =

iFROM-pred-insert-conv

iIN-pred-insert-conv

iMOD-pred-insert-conv

iMODb-pred-insert-conv

lemmas *iT-Suc-pred-insert* =

iFROM-Suc-pred-insert-conv

iIN-Suc-pred-insert-conv

iMOD-Suc-pred-insert-conv

iMODb-Suc-pred-insert-conv

lemma *iMOD-mem-diff*: $\llbracket a \in [r, \text{ mod } m]; b \in [r, \text{ mod } m] \rrbracket \implies (a - b) \text{ mod } m = 0$

<proof>

lemma *iMODb-mem-diff*: $\llbracket a \in [r, \text{ mod } m, c]; b \in [r, \text{ mod } m, c] \rrbracket \implies (a - b) \text{ mod } m = 0$

<proof>

1.1.3 Interval conversions

lemma *iIN-0-iTILL-conv*: $[0 \dots, n] = [\dots n]$

<proof>

lemma *iIN-iTILL-iTILL-conv*: $0 < n \implies [n \dots, d] = [\dots n + d] - [\dots n - \text{Suc } 0]$

<proof>

lemma *iIN-iFROM-iTILL-conv*: $[n \dots, d] = [n \dots] \cap [\dots n + d]$

<proof>

lemma *iMODb-iMOD-iTILL-conv*: $[r, \text{ mod } m, c] = [r, \text{ mod } m] \cap [\dots r + m * c]$

<proof>

lemma *iMODb-iMOD-iIN-conv*: $[r, \text{ mod } m, c] = [r, \text{ mod } m] \cap [r \dots, m * c]$

<proof>

lemma *iFROM-iTILL-iIN-conv*: $n \leq n' \implies [n \dots] \cap [\dots n'] = [n \dots, n' - n]$

<proof>

lemma *iMOD-iTILL-iMODb-conv*:

$$r \leq n \implies [r, \text{mod } m] \cap [\dots n] = [r, \text{mod } m, (n - r) \text{ div } m]$$

<proof>

lemma *iMOD-iIN-iMODb-conv*:

$$[r, \text{mod } m] \cap [r\dots, d] = [r, \text{mod } m, d \text{ div } m]$$

<proof>

lemma *iFROM-0*: $[0\dots] = \text{UNIV}$

<proof>

lemma *iTILL-0*: $[\dots 0] = \{0\}$

<proof>

lemma *iIN-0*: $[n\dots, 0] = \{n\}$

<proof>

lemma *iMOD-0*: $[r, \text{mod } 0] = [r\dots, 0]$

<proof>

lemma *iMODb-mod-0*: $[r, \text{mod } 0, c] = [r\dots, 0]$

<proof>

lemma *iMODb-0*: $[r, \text{mod } m, 0] = [r\dots, 0]$

<proof>

lemmas *iT-0* =

iFROM-0

iTILL-0

iIN-0

iMOD-0

iMODb-mod-0

iMODb-0

lemma *iMOD-1*: $[r, \text{mod } \text{Suc } 0] = [r\dots]$

<proof>

lemma *iMODb-mod-1*: $[r, \text{mod } \text{Suc } 0, c] = [r\dots, c]$

<proof>

1.1.4 Finiteness and emptiness of intervals

lemma

iFROM-not-empty: $[n\dots] \neq \{\}$ **and**

iTILL-not-empty: $[\dots n] \neq \{\}$ **and**

iIN-not-empty: $[n\dots, d] \neq \{\}$ **and**

iMOD-not-empty: $[r, \text{mod } m] \neq \{\}$ **and**

iMODb-not-empty: $[r, \text{mod } m, c] \neq \{\}$

<proof>

lemmas *iT-not-empty* =
iFROM-not-empty
iTILL-not-empty
iIN-not-empty
iMOD-not-empty
iMODb-not-empty

lemma
iTILL-finite: *finite* [$\dots n$] **and**
iIN-finite: *finite* [$n \dots, d$] **and**
iMODb-finite: *finite* [$r, \text{mod } m, c$] **and**
iMOD-0-finite: *finite* [$r, \text{mod } 0$]
⟨*proof*⟩

lemma *iFROM-infinite*: *infinite* [$n \dots$]
⟨*proof*⟩

lemma *iMOD-infinite*: $0 < m \implies \textit{infinite}$ [$r, \text{mod } m$]
⟨*proof*⟩

lemmas *iT-finite* =
iTILL-finite
iIN-finite
iMODb-finite *iMOD-0-finite*

lemmas *iT-infinite* =
iFROM-infinite
iMOD-infinite

1.1.5 *Min* and *Max* element of an interval

lemma
iTILL-Min: *iMin* [$\dots n$] = 0 **and**
iFROM-Min: *iMin* [$n \dots$] = n **and**
iIN-Min: *iMin* [$n \dots, d$] = n **and**
iMOD-Min: *iMin* [$r, \text{mod } m$] = r **and**
iMODb-Min: *iMin* [$r, \text{mod } m, c$] = r
⟨*proof*⟩

lemmas *iT-Min* =
iIN-Min
iTILL-Min
iFROM-Min
iMOD-Min
iMODb-Min

lemma
iTILL-Max: *Max* [$\dots n$] = n **and**

iIN-Max: $\text{Max } [n..,d] = n+d$ **and**
iMODb-Max: $\text{Max } [r, \text{mod } m, c] = r + m * c$ **and**
iMOD-0-Max: $\text{Max } [r, \text{mod } 0] = r$
 ⟨*proof*⟩

lemmas *iT-Max* =
iTILL-Max
iIN-Max
iMODb-Max
iMOD-0-Max

lemma
iTILL-iMax: $i\text{Max } [..n] = \text{enat } n$ **and**
iIN-iMax: $i\text{Max } [n..,d] = \text{enat } (n+d)$ **and**
iMODb-iMax: $i\text{Max } [r, \text{mod } m, c] = \text{enat } (r + m * c)$ **and**
iMOD-0-iMax: $i\text{Max } [r, \text{mod } 0] = \text{enat } r$ **and**
iFROM-iMax: $i\text{Max } [n..] = \infty$ **and**
iMOD-iMax: $0 < m \implies i\text{Max } [r, \text{mod } m] = \infty$
 ⟨*proof*⟩

lemmas *iT-iMax* =
iTILL-iMax
iIN-iMax
iMODb-iMax
iMOD-0-iMax
iFROM-iMax
iMOD-iMax

1.2 Adding and subtracting constants to interval elements

lemma
iFROM-plus: $x \in [n..] \implies x + k \in [n..]$ **and**
iFROM-Suc: $x \in [n..] \implies \text{Suc } x \in [n..]$ **and**
iFROM-minus: $\llbracket x \in [n..]; k \leq x - n \rrbracket \implies x - k \in [n..]$ **and**
iFROM-pred: $n < x \implies x - \text{Suc } 0 \in [n..]$
 ⟨*proof*⟩

lemma
iTILL-plus: $\llbracket x \in [..n]; k \leq n - x \rrbracket \implies x + k \in [..n]$ **and**
iTILL-Suc: $x < n \implies \text{Suc } x \in [..n]$ **and**
iTILL-minus: $x \in [..n] \implies x - k \in [..n]$ **and**
iTILL-pred: $x \in [..n] \implies x - \text{Suc } 0 \in [..n]$
 ⟨*proof*⟩

lemma *iIN-plus*: $\llbracket x \in [n..,d]; k \leq n + d - x \rrbracket \implies x + k \in [n..,d]$
 ⟨*proof*⟩

lemma *iIN-Suc*: $\llbracket x \in [n..,d]; x < n + d \rrbracket \implies \text{Suc } x \in [n..,d]$
 ⟨*proof*⟩

lemma *iIN-minus*: $\llbracket x \in [n..d]; k \leq x - n \rrbracket \implies x - k \in [n..d]$
 ⟨proof⟩

lemma *iIN-pred*: $\llbracket x \in [n..d]; n < x \rrbracket \implies x - \text{Suc } 0 \in [n..d]$
 ⟨proof⟩

lemma *iMOD-plus-divisor-mult*: $x \in [r, \text{mod } m] \implies x + k * m \in [r, \text{mod } m]$
 ⟨proof⟩

corollary *iMOD-plus-divisor*: $x \in [r, \text{mod } m] \implies x + m \in [r, \text{mod } m]$
 ⟨proof⟩

lemma *iMOD-minus-divisor-mult*:
 $\llbracket x \in [r, \text{mod } m]; k * m \leq x - r \rrbracket \implies x - k * m \in [r, \text{mod } m]$
 ⟨proof⟩

corollary *iMOD-minus-divisor-mult2*:
 $\llbracket x \in [r, \text{mod } m]; k \leq (x - r) \text{ div } m \rrbracket \implies x - k * m \in [r, \text{mod } m]$
 ⟨proof⟩

corollary *iMOD-minus-divisor*:
 $\llbracket x \in [r, \text{mod } m]; m + r \leq x \rrbracket \implies x - m \in [r, \text{mod } m]$
 ⟨proof⟩

lemma *iMOD-plus*:
 $x \in [r, \text{mod } m] \implies (x + k \in [r, \text{mod } m]) = (k \text{ mod } m = 0)$
 ⟨proof⟩

corollary *iMOD-Suc*:
 $x \in [r, \text{mod } m] \implies (\text{Suc } x \in [r, \text{mod } m]) = (m = \text{Suc } 0)$
 ⟨proof⟩

lemma *iMOD-minus*:
 $\llbracket x \in [r, \text{mod } m]; k \leq x - r \rrbracket \implies (x - k \in [r, \text{mod } m]) = (k \text{ mod } m = 0)$
 ⟨proof⟩

corollary *iMOD-pred*:
 $\llbracket x \in [r, \text{mod } m]; r < x \rrbracket \implies (x - \text{Suc } 0 \in [r, \text{mod } m]) = (m = \text{Suc } 0)$
 ⟨proof⟩

lemma *iMODb-plus-divisor-mult*:
 $\llbracket x \in [r, \text{mod } m, c]; k * m \leq r + m * c - x \rrbracket \implies x + k * m \in [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-plus-divisor-mult2*:
 $\llbracket x \in [r, \text{mod } m, c]; k \leq c - (x - r) \text{ div } m \rrbracket \implies$
 $x + k * m \in [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-plus-divisor*:

$\llbracket x \in [r, \text{mod } m, c]; x < r + m * c \rrbracket \implies x + m \in [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-minus-divisor-mult*:

$\llbracket x \in [r, \text{mod } m, c]; r + k * m \leq x \rrbracket \implies x - k * m \in [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-plus*:

$\llbracket x \in [r, \text{mod } m, c]; k \leq r + m * c - x \rrbracket \implies$
 $(x + k \in [r, \text{mod } m, c]) = (k \text{ mod } m = 0)$
 ⟨proof⟩

corollary *iMODb-Suc*:

$\llbracket x \in [r, \text{mod } m, c]; x < r + m * c \rrbracket \implies$
 $(\text{Suc } x \in [r, \text{mod } m, c]) = (m = \text{Suc } 0)$
 ⟨proof⟩

lemma *iMODb-minus*:

$\llbracket x \in [r, \text{mod } m, c]; k \leq x - r \rrbracket \implies$
 $(x - k \in [r, \text{mod } m, c]) = (k \text{ mod } m = 0)$
 ⟨proof⟩

corollary *iMODb-pred*:

$\llbracket x \in [r, \text{mod } m, c]; r < x \rrbracket \implies$
 $(x - \text{Suc } 0 \in [r, \text{mod } m, c]) = (m = \text{Suc } 0)$
 ⟨proof⟩

lemmas *iFROM-plus-minus =*

iFROM-plus
iFROM-Suc
iFROM-minus
iFROM-pred

lemmas *iTILL-plus-minus =*

iTILL-plus
iTILL-Suc
iTILL-minus
iTILL-pred

lemmas *iIN-plus-minus =*

iIN-plus
iIN-Suc
iTILL-minus
iIN-pred

lemmas *iMOD-plus-minus-divisor =*

iMOD-plus-divisor-mult
iMOD-plus-divisor
iMOD-minus-divisor-mult

iMOD-minus-divisor-mult2
iMOD-minus-divisor

lemmas *iMOD-plus-minus* =
iMOD-plus
iMOD-Suc
iMOD-minus
iMOD-pred

lemmas *iMODb-plus-minus-divisor* =
iMODb-plus-divisor-mult
iMODb-plus-divisor-mult2
iMODb-plus-divisor
iMODb-minus-divisor-mult

lemmas *iMODb-plus-minus* =
iMODb-plus
iMODb-Suc
iMODb-minus
iMODb-pred

lemmas *iT-plus-minus* =
iFROM-plus-minus
iTILL-plus-minus
iIN-plus-minus
iMOD-plus-minus-divisor
iMOD-plus-minus
iMODb-plus-minus-divisor
iMODb-plus-minus

1.3 Relations between intervals

1.3.1 Auxiliary lemmata

lemma *Suc-in-imp-not-subset-iMOD*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies \neg S \subseteq [r, \text{mod } m]$
 ⟨proof⟩

corollary *Suc-in-imp-neq-iMOD*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies S \neq [r, \text{mod } m]$
 ⟨proof⟩

lemma *Suc-in-imp-not-subset-iMODb*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies \neg S \subseteq [r, \text{mod } m, c]$
 ⟨proof⟩

corollary *Suc-in-imp-neq-iMODb*:

$\llbracket n \in S; \text{Suc } n \in S; m \neq \text{Suc } 0 \rrbracket \implies S \neq [r, \text{mod } m, c]$
 ⟨proof⟩

1.3.2 Subset relation between intervals

lemma

iIN-iFROM-subset-same: $[n\dots,d] \subseteq [n\dots]$ **and**
iIN-iTILL-subset-same: $[n\dots,d] \subseteq [\dots,n+d]$ **and**
iMOD-iFROM-subset-same: $[r, \text{mod } m] \subseteq [r\dots]$ **and**
iMODb-iTILL-subset-same: $[r, \text{mod } m, c] \subseteq [\dots,r+m*c]$ **and**
iMODb-iIN-subset-same: $[r, \text{mod } m, c] \subseteq [r\dots,m*c]$ **and**
iMODb-iMOD-subset-same: $[r, \text{mod } m, c] \subseteq [r, \text{mod } m]$

<proof>

lemmas *iT-subset-same* =

iIN-iFROM-subset-same
iIN-iTILL-subset-same
iMOD-iFROM-subset-same
iMODb-iTILL-subset-same
iMODb-iIN-subset-same
iMODb-iTILL-subset-same
iMODb-iMOD-subset-same

lemma *iMODb-imp-iMOD*: $x \in [r, \text{mod } m, c] \implies x \in [r, \text{mod } m]$

<proof>

lemma *iMOD-imp-iMODb*:

$\llbracket x \in [r, \text{mod } m]; x \leq r + m * c \rrbracket \implies x \in [r, \text{mod } m, c]$

<proof>

lemma *iMOD-singleton-subset-conv*: $([r, \text{mod } m] \subseteq \{a\}) = (r = a \wedge m = 0)$

<proof>

lemma *iMOD-singleton-eq-conv*: $([r, \text{mod } m] = \{a\}) = (r = a \wedge m = 0)$

<proof>

lemma *iMODb-singleton-subset-conv*:

$([r, \text{mod } m, c] \subseteq \{a\}) = (r = a \wedge (m = 0 \vee c = 0))$

<proof>

lemma *iMODb-singleton-eq-conv*:

$([r, \text{mod } m, c] = \{a\}) = (r = a \wedge (m = 0 \vee c = 0))$

<proof>

lemma *iMODb-subset-imp-divisor-mod-0*:

$\llbracket 0 < c'; [r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c] \rrbracket \implies m' \text{ mod } m = 0$

<proof>

lemma *iMOD-subset-imp-divisor-mod-0*:

$[r', \text{mod } m'] \subseteq [r, \text{mod } m] \implies m' \text{ mod } m = 0$

<proof>

lemma *iMOD-subset-imp-iMODb-subset*:

$\llbracket [r', \text{mod } m'] \subseteq [r, \text{mod } m]; r' + m' * c' \leq r + m * c \rrbracket \implies [r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]$

$\langle proof \rangle$

lemma *iMODb-subset-imp-iMOD-subset*:

$$\llbracket [r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]; 0 < c' \rrbracket \implies \\ [r', \text{mod } m'] \subseteq [r, \text{mod } m]$$

$\langle proof \rangle$

lemma *iMODb-0-iMOD-subset-conv*:

$$([r', \text{mod } m', 0] \subseteq [r, \text{mod } m]) = \\ (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r')$$

$\langle proof \rangle$

lemma *iFROM-subset-conv*: $([n' \dots] \subseteq [n \dots]) = (n \leq n')$

$\langle proof \rangle$

lemma *iFROM-iMOD-subset-conv*: $([n' \dots] \subseteq [r, \text{mod } m]) = (r \leq n' \wedge m = \text{Suc } 0)$

$\langle proof \rangle$

lemma *iIN-subset-conv*: $([n' \dots, d'] \subseteq [n \dots, d]) = (n \leq n' \wedge n' + d' \leq n + d)$

$\langle proof \rangle$

lemma *iIN-iFROM-subset-conv*: $([n' \dots, d'] \subseteq [n \dots]) = (n \leq n')$

$\langle proof \rangle$

lemma *iIN-iTILL-subset-conv*: $([n' \dots, d'] \subseteq [\dots n]) = (n' + d' \leq n)$

$\langle proof \rangle$

lemma *iIN-iMOD-subset-conv*:

$$0 < d' \implies ([n' \dots, d'] \subseteq [r, \text{mod } m]) = (r \leq n' \wedge m = \text{Suc } 0)$$

$\langle proof \rangle$

lemma *iIN-iMODb-subset-conv*:

$$0 < d' \implies \\ ([n' \dots, d'] \subseteq [r, \text{mod } m, c]) = \\ (r \leq n' \wedge m = \text{Suc } 0 \wedge n' + d' \leq r + m * c)$$

$\langle proof \rangle$

lemma *iTILL-subset-conv*: $([\dots n'] \subseteq [\dots n]) = (n' \leq n)$

$\langle proof \rangle$

lemma *iTILL-iFROM-subset-conv*: $([\dots n'] \subseteq [n \dots]) = (n = 0)$

$\langle proof \rangle$

lemma *iTILL-iIN-subset-conv*: $([\dots n'] \subseteq [n \dots, d]) = (n = 0 \wedge n' \leq d)$

$\langle proof \rangle$

lemma *iTILL-iMOD-subset-conv*:

$$0 < n' \implies ([..n'] \subseteq [r, \text{mod } m]) = (r = 0 \wedge m = \text{Suc } 0)$$

<proof>

lemma *iTILL-iMODb-subset-conv:*

$$0 < n' \implies ([..n'] \subseteq [r, \text{mod } m, c]) = (r = 0 \wedge m = \text{Suc } 0 \wedge n' \leq r + m * c)$$

<proof>

lemma *iMOD-iFROM-subset-conv:* $([r', \text{mod } m'] \subseteq [n..]) = (n \leq r')$

<proof>

lemma *iMODb-iFROM-subset-conv:* $([r', \text{mod } m', c'] \subseteq [n..]) = (n \leq r')$

<proof>

lemma *iMODb-iIN-subset-conv:*

$$([r', \text{mod } m', c'] \subseteq [n..,d]) = (n \leq r' \wedge r' + m' * c' \leq n + d)$$

<proof>

lemma *iMODb-iTILL-subset-conv:*

$$([r', \text{mod } m', c'] \subseteq [..n]) = (r' + m' * c' \leq n)$$

<proof>

lemma *iMOD-0-subset-conv:* $([r', \text{mod } 0] \subseteq [r, \text{mod } m]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r')$

<proof>

lemma *iMOD-subset-conv:* $0 < m \implies$

$$([r', \text{mod } m'] \subseteq [r, \text{mod } m]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0)$$

<proof>

lemma *iMODb-subset-mod-0-conv:*

$$([r', \text{mod } m', c'] \subseteq [r, \text{mod } 0, c]) = (r'=r \wedge (m'=0 \vee c'=0))$$

<proof>

lemma *iMODb-subset-0-conv:*

$$([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, 0]) = (r'=r \wedge (m'=0 \vee c'=0))$$

<proof>

lemma *iMODb-0-subset-conv:*

$$([r', \text{mod } m', 0] \subseteq [r, \text{mod } m, c]) = (r' \in [r, \text{mod } m, c])$$

<proof>

lemma *iMODb-mod-0-subset-conv:*

$$([r', \text{mod } 0, c'] \subseteq [r, \text{mod } m, c]) = (r' \in [r, \text{mod } m, c])$$

<proof>

lemma *iMODb-subset-conv':* $\llbracket 0 < c; 0 < c' \rrbracket \implies$

$$([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]) = (r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0 \wedge$$

$r' + m' * c' \leq r + m * c$
 ⟨proof⟩

lemma *iMODb-subset-conv*: $\llbracket 0 < m'; 0 < c' \rrbracket \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0 \wedge$
 $r' + m' * c' \leq r + m * c)$
 ⟨proof⟩

lemma *iMODb-iMOD-subset-conv*: $0 < c' \implies$
 $([r', \text{mod } m', c'] \subseteq [r, \text{mod } m]) =$
 $(r' \text{ mod } m = r \text{ mod } m \wedge r \leq r' \wedge m' \text{ mod } m = 0)$
 ⟨proof⟩

lemmas *iT-subset-conv* =
iFROM-subset-conv
iFROM-iMOD-subset-conv
iTILL-subset-conv
iTILL-iFROM-subset-conv
iTILL-iIN-subset-conv
iTILL-iMOD-subset-conv
iTILL-iMODb-subset-conv
iIN-subset-conv
iIN-iFROM-subset-conv
iIN-iTILL-subset-conv
iIN-iMOD-subset-conv
iIN-iMODb-subset-conv
iMOD-subset-conv
iMOD-iFROM-subset-conv
iMODb-subset-conv'
iMODb-subset-conv
iMODb-iFROM-subset-conv
iMODb-iIN-subset-conv
iMODb-iTILL-subset-conv
iMODb-iMOD-subset-conv

lemma *iFROM-subset*: $n \leq n' \implies [n'..] \subseteq [n..]$
 ⟨proof⟩

lemma *not-iFROM-iIN-subset*: $\neg [n'..] \subseteq [n.., d]$
 ⟨proof⟩

lemma *not-iFROM-iTILL-subset*: $\neg [n'..] \subseteq [...n]$
 ⟨proof⟩

lemma *not-iFROM-iMOD-subset*: $m \neq \text{Suc } 0 \implies \neg [n'..] \subseteq [r, \text{mod } m]$
 ⟨proof⟩

lemma *not-iFROM-iMODb-subset*: $\neg [n'..] \subseteq [r, \text{mod } m, c]$

<proof>

lemma *iIN-subset*: $\llbracket n \leq n'; n' + d' \leq n + d \rrbracket \implies [n' \dots, d'] \subseteq [n \dots, d]$
<proof>

lemma *iIN-iFROM-subset*: $n \leq n' \implies [n' \dots, d'] \subseteq [n \dots]$
<proof>

lemma *iIN-iTILL-subset*: $n' + d' \leq n \implies [n' \dots, d'] \subseteq [\dots n]$
<proof>

lemma *not-iIN-iMODb-subset*: $\llbracket 0 < d'; m \neq \text{Suc } 0 \rrbracket \implies \neg [n' \dots, d'] \subseteq [r, \text{mod } m, c]$
<proof>

lemma *not-iIN-iMOD-subset*: $\llbracket 0 < d'; m \neq \text{Suc } 0 \rrbracket \implies \neg [n' \dots, d'] \subseteq [r, \text{mod } m]$
<proof>

lemma *iTILL-subset*: $n' \leq n \implies [\dots n'] \subseteq [\dots n]$
<proof>

lemma *iTILL-iFROM-subset*: $([\dots n'] \subseteq [0 \dots])$
<proof>

lemma *iTILL-iIN-subset*: $n' \leq d \implies ([\dots n'] \subseteq [0 \dots, d])$
<proof>

lemma *not-iTILL-iMOD-subset*:
 $\llbracket 0 < n'; m \neq \text{Suc } 0 \rrbracket \implies \neg [\dots n'] \subseteq [r, \text{mod } m]$
<proof>

lemma *not-iTILL-iMODb-subset*:
 $\llbracket 0 < n'; m \neq \text{Suc } 0 \rrbracket \implies \neg [\dots n'] \subseteq [r, \text{mod } m, c]$
<proof>

lemma *iMOD-iFROM-subset*: $n \leq r' \implies [r', \text{mod } m'] \subseteq [n \dots]$
<proof>

lemma *not-iMOD-iIN-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [n \dots, d]$
<proof>

lemma *not-iMOD-iTILL-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [\dots n]$
<proof>

lemma *iMOD-subset*:
 $\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0 \rrbracket \implies [r', \text{mod } m'] \subseteq [r, \text{mod } m]$
<proof>

lemma *not-iMOD-iMODb-subset*: $0 < m' \implies \neg [r', \text{mod } m'] \subseteq [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-iFROM-subset*: $n \leq r' \implies [r', \text{mod } m', c'] \subseteq [n..]$
 ⟨proof⟩

lemma *iMODb-iTILL-subset*:
 $r' + m' * c' \leq n \implies [r', \text{mod } m', c'] \subseteq [..n]$
 ⟨proof⟩

lemma *iMODb-iIN-subset*:
 $\llbracket n \leq r'; r' + m' * c' \leq n + d \rrbracket \implies [r', \text{mod } m', c'] \subseteq [n..,d]$
 ⟨proof⟩

lemma *iMODb-iMOD-subset*:
 $\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0 \rrbracket \implies [r', \text{mod } m', c'] \subseteq [r, \text{mod } m]$
 ⟨proof⟩

lemma *iMODb-subset*:
 $\llbracket r \leq r'; r' \text{ mod } m = r \text{ mod } m; m' \text{ mod } m = 0; r' + m' * c' \leq r + m * c \rrbracket \implies$
 $[r', \text{mod } m', c'] \subseteq [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iFROM-trans*: $\llbracket y \in [x..]; z \in [y..] \rrbracket \implies z \in [x..]$
 ⟨proof⟩

lemma *iTILL-trans*: $\llbracket y \in [..x]; z \in [..y] \rrbracket \implies z \in [..x]$
 ⟨proof⟩

lemma *iIN-trans*:
 $\llbracket y \in [x..,d]; z \in [y..,d']; d' \leq x + d - y \rrbracket \implies z \in [x..,d]$
 ⟨proof⟩

lemma *iMOD-trans*:
 $\llbracket y \in [x, \text{mod } m]; z \in [y, \text{mod } m] \rrbracket \implies z \in [x, \text{mod } m]$
 ⟨proof⟩

lemma *iMODb-trans*:
 $\llbracket y \in [x, \text{mod } m, c]; z \in [y, \text{mod } m, c']; m * c' \leq x + m * c - y \rrbracket \implies$
 $z \in [x, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-trans'*:
 $\llbracket y \in [x, \text{mod } m, c]; z \in [y, \text{mod } m, c']; c' \leq x \text{ div } m + c - y \text{ div } m \rrbracket \implies$
 $z \in [x, \text{mod } m, c]$
 ⟨proof⟩

1.3.3 Equality of intervals

lemma *iFROM-eq-conv*: $([n\dots] = [n'\dots]) = (n = n')$
 ⟨proof⟩

lemma *iIN-eq-conv*: $([n\dots, d] = [n'\dots, d']) = (n = n' \wedge d = d')$
 ⟨proof⟩

lemma *iTILL-eq-conv*: $([\dots n] = [\dots n']) = (n = n')$
 ⟨proof⟩

lemma *iMOD-0-eq-conv*: $([r, \text{mod } 0] = [r', \text{mod } m']) = (r = r' \wedge m' = 0)$
 ⟨proof⟩

lemma *iMOD-eq-conv*: $0 < m \implies ([r, \text{mod } m] = [r', \text{mod } m']) = (r = r' \wedge m = m')$
 ⟨proof⟩

lemma *iMODb-mod-0-eq-conv*:
 $([r, \text{mod } 0, c] = [r', \text{mod } m', c']) = (r = r' \wedge (m' = 0 \vee c' = 0))$
 ⟨proof⟩

lemma *iMODb-0-eq-conv*:
 $([r, \text{mod } m, 0] = [r', \text{mod } m', c']) = (r = r' \wedge (m' = 0 \vee c' = 0))$
 ⟨proof⟩

lemma *iMODb-eq-conv*: $\llbracket 0 < m; 0 < c \rrbracket \implies$
 $([r, \text{mod } m, c] = [r', \text{mod } m', c']) = (r = r' \wedge m = m' \wedge c = c')$
 ⟨proof⟩

lemma *iMOD-iFROM-eq-conv*: $([n\dots] = [r, \text{mod } m]) = (n = r \wedge m = \text{Suc } 0)$
 ⟨proof⟩

lemma *iMODb-iIN-0-eq-conv*:
 $([n\dots, 0] = [r, \text{mod } m, c]) = (n = r \wedge (m = 0 \vee c = 0))$
 ⟨proof⟩

lemma *iMODb-iIN-eq-conv*:
 $0 < d \implies ([n\dots, d] = [r, \text{mod } m, c]) = (n = r \wedge m = \text{Suc } 0 \wedge c = d)$
 ⟨proof⟩

1.3.4 Inequality of intervals

lemma *iFROM-iIN-neg*: $[n'\dots] \neq [n\dots, d]$
 ⟨proof⟩

corollary *iFROM-iTILL-neg*: $[n'\dots] \neq [\dots n]$
 ⟨proof⟩

corollary *iFROM-iMOD-neg*: $m \neq \text{Suc } 0 \implies [n\dots] \neq [r, \text{mod } m]$

<proof>

corollary *iFROM-iMODb-neq*: $[n\dots] \neq [r, \text{mod } m, c]$

<proof>

corollary *iIN-iMOD-neq*: $0 < m \implies [n\dots, d] \neq [r, \text{mod } m]$

<proof>

corollary *iIN-iMODb-neq2*: $\llbracket m \neq \text{Suc } 0; 0 < d \rrbracket \implies [n\dots, d] \neq [r, \text{mod } m, c]$

<proof>

lemma *iIN-iMODb-neq*: $\llbracket 2 \leq m; 0 < c \rrbracket \implies [n\dots, d] \neq [r, \text{mod } m, c]$

<proof>

lemma *iTILL-iIN-neq*: $0 < n \implies [\dots n^\wedge] \neq [n\dots, d]$

<proof>

corollary *iTILL-iMOD-neq*: $0 < m \implies [\dots n] \neq [r, \text{mod } m]$

<proof>

corollary *iTILL-iMODb-neq*:

$\llbracket m \neq \text{Suc } 0; 0 < n \rrbracket \implies [\dots n] \neq [r, \text{mod } m, c]$

<proof>

lemma *iMOD-iMODb-neq*: $0 < m \implies [r, \text{mod } m] \neq [r', \text{mod } m', c^\wedge]$

<proof>

lemmas *iT-neq* =

iFROM-iTILL-neq *iFROM-iIN-neq* *iFROM-iMOD-neq* *iFROM-iMODb-neq*

iTILL-iIN-neq *iTILL-iMOD-neq* *iTILL-iMODb-neq*

iIN-iMOD-neq *iIN-iMODb-neq* *iIN-iMODb-neq2*

iMOD-iMODb-neq

1.4 Union and intersection of intervals

lemma *iFROM-union'*: $[n\dots] \cup [n'\dots] = [\min n n'\dots]$

<proof>

corollary *iFROM-union*: $n \leq n' \implies [n\dots] \cup [n'\dots] = [n\dots]$

<proof>

lemma *iFROM-inter'*: $[n\dots] \cap [n'\dots] = [\max n n'\dots]$

<proof>

corollary *iFROM-inter*: $n' \leq n \implies [n\dots] \cap [n'\dots] = [n\dots]$

<proof>

lemma *iTILL-union'*: $[\dots n] \cup [\dots n^\wedge] = [\dots \max n n^\wedge]$

<proof>

corollary *iTILL-union*: $n' \leq n \implies [\dots n] \cup [\dots n'] = [\dots n]$
 ⟨proof⟩

lemma *iTILL-iFROM-union*: $n \leq n' \implies [\dots n'] \cup [n\dots] = UNIV$
 ⟨proof⟩

lemma *iTILL-inter'*: $[\dots n] \cap [\dots n'] = [\dots \min n n']$
 ⟨proof⟩

corollary *iTILL-inter*: $n \leq n' \implies [\dots n] \cap [\dots n'] = [\dots n]$
 ⟨proof⟩

Union and intersection for iIN only when there intersection is not empty and none of them is other's subset, for instance: .. n .. n+d .. n' .. n'+d'

lemma *iIN-union*:
 $\llbracket n \leq n'; n' \leq \text{Suc } (n + d); n + d \leq n' + d' \rrbracket \implies$
 $[n\dots, d] \cup [n'\dots, d'] = [n\dots, n' - n + d']$
 ⟨proof⟩

lemma *iIN-append-union*:
 $[n\dots, d] \cup [n + d\dots, d'] = [n\dots, d + d']$
 ⟨proof⟩

lemma *iIN-append-union-Suc*:
 $[n\dots, d] \cup [\text{Suc } (n + d)\dots, d'] = [n\dots, \text{Suc } (d + d')]$
 ⟨proof⟩

lemma *iIN-append-union-pred*:
 $0 < d \implies [n\dots, d - \text{Suc } 0] \cup [n + d\dots, d'] = [n\dots, d + d']$
 ⟨proof⟩

lemma *iIN-iFROM-union*:
 $n' \leq \text{Suc } (n + d) \implies [n\dots, d] \cup [n'\dots] = [\min n n'\dots]$
 ⟨proof⟩

lemma *iIN-iFROM-append-union*:
 $[n\dots, d] \cup [n + d\dots] = [n\dots]$
 ⟨proof⟩

lemma *iIN-iFROM-append-union-Suc*:
 $[n\dots, d] \cup [\text{Suc } (n + d)\dots] = [n\dots]$
 ⟨proof⟩

lemma *iIN-iFROM-append-union-pred*:
 $0 < d \implies [n\dots, d - \text{Suc } 0] \cup [n + d\dots] = [n\dots]$
 ⟨proof⟩

lemma *iN-inter*:

$$\begin{aligned} & \llbracket n \leq n'; n' \leq n + d; n + d \leq n' + d' \rrbracket \implies \\ & [n \dots, d] \cap [n' \dots, d'] = [n' \dots, n + d - n'] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m \rrbracket \implies \\ & [r, \bmod m] \cup [r', \bmod m] = [r, \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-union'*:

$$\begin{aligned} & r \bmod m = r' \bmod m \implies \\ & [r, \bmod m] \cup [r', \bmod m] = [\min r r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-inter*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m \rrbracket \implies \\ & [r, \bmod m] \cap [r', \bmod m] = [r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMOD-inter'*:

$$\begin{aligned} & r \bmod m = r' \bmod m \implies \\ & [r, \bmod m] \cap [r', \bmod m] = [\max r r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * c; r + m * c \leq r' + m * c' \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m, c'] = [r, \bmod m, r' \text{ div } m - r \text{ div } m + c'] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-append-union*:

$$\begin{aligned} & [r, \bmod m, c] \cup [r + m * c, \bmod m, c'] = [r, \bmod m, c + c'] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-iMOD-append-union'*:

$$\begin{aligned} & \llbracket r \bmod m = r' \bmod m; r' \leq r + m * \text{Suc } c \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m] = [\min r r', \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-iMOD-append-union*:

$$\begin{aligned} & \llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * \text{Suc } c \rrbracket \implies \\ & [r, \bmod m, c] \cup [r', \bmod m] = [r, \bmod m] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-append-union-Suc*:

$$\begin{aligned} & [r, \bmod m, c] \cup [r + m * \text{Suc } c, \bmod m, c'] = [r, \bmod m, \text{Suc } (c + c')] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *iMODb-append-union-pred*:

$$0 < c \implies [r, \bmod m, c - \text{Suc } 0] \cup [r + m * c, \bmod m, c'] = [r, \bmod m, c +$$

c'
 $\langle proof \rangle$

lemma *iMODb-inter*:

$\llbracket r \leq r'; r \bmod m = r' \bmod m; r' \leq r + m * c; r + m * c \leq r' + m * c' \rrbracket \implies$
 $[r, \bmod m, c] \cap [r', \bmod m, c'] = [r', \bmod m, c - (r'-r) \text{ div } m]$
 $\langle proof \rangle$

lemmas *iT-union'* =

iFROM-union'
iTILL-union'
iMOD-union'
iMODb-iMOD-append-union'

lemmas *iT-union* =

iFROM-union
iTILL-union
iTILL-iFROM-union
iIN-union
iIN-iFROM-union
iMOD-union
iMODb-union

lemmas *iT-union-append* =

iIN-append-union
iIN-append-union-Suc
iIN-append-union-pred
iIN-iFROM-append-union
iIN-iFROM-append-union-Suc
iIN-iFROM-append-union-pred
iMODb-append-union
iMODb-iMOD-append-union
iMODb-append-union-Suc
iMODb-append-union-pred

lemmas *iT-inter'* =

iFROM-inter'
iTILL-inter'
iMOD-inter'

lemmas *iT-inter* =

iFROM-inter
iTILL-inter
iIN-inter
iMOD-inter
iMODb-inter

lemma *mod-partition-Union*:

$0 < m \implies (\bigcup k. A \cap [k * m .. m - \text{Suc } 0]) = A$

$\langle \text{proof} \rangle$

lemma *finite-mod-partition-Union*:

$\llbracket 0 < m; \text{finite } A \rrbracket \implies$
 $(\bigcup_{k \leq \text{Max } A \text{ div } m} A \cap [k * m \dots, m - \text{Suc } 0]) = A$
 $\langle \text{proof} \rangle$

lemma *mod-partition-is-disjoint*:

$\llbracket 0 < (m :: \text{nat}); k \neq k' \rrbracket \implies$
 $(A \cap [k * m \dots, m - \text{Suc } 0]) \cap (A \cap [k' * m \dots, m - \text{Suc } 0]) = \{\}$
 $\langle \text{proof} \rangle$

1.5 Cutting intervals

lemma *iTILL-cut-le*: $[\dots n] \downarrow \leq t = (\text{if } t \leq n \text{ then } [\dots t] \text{ else } [\dots n])$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-le1*: $t \in [\dots n] \implies [\dots n] \downarrow \leq t = [\dots t]$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-le2*: $t \notin [\dots n] \implies [\dots n] \downarrow \leq t = [\dots n]$
 $\langle \text{proof} \rangle$

lemma *iFROM-cut-le*:

$[n \dots] \downarrow \leq t =$
 $(\text{if } t < n \text{ then } \{\} \text{ else } [n \dots, t - n])$
 $\langle \text{proof} \rangle$

corollary *iFROM-cut-le1*: $t \in [n \dots] \implies [n \dots] \downarrow \leq t = [n \dots, t - n]$
 $\langle \text{proof} \rangle$

lemma *iIN-cut-le*:

$[n \dots, d] \downarrow \leq t =$
 $(\text{if } t < n \text{ then } \{\} \text{ else}$
 $\text{if } t \leq n + d \text{ then } [n \dots, t - n]$
 $\text{else } [n \dots, d])$
 $\langle \text{proof} \rangle$

corollary *iIN-cut-le1*:

$t \in [n \dots, d] \implies [n \dots, d] \downarrow \leq t = [n \dots, t - n]$
 $\langle \text{proof} \rangle$

lemma *iMOD-cut-le*:

$[r, \text{mod } m] \downarrow \leq t =$
 $(\text{if } t < r \text{ then } \{\}$
 $\text{else } [r, \text{mod } m, (t - r) \text{ div } m])$
 $\langle \text{proof} \rangle$

lemma *iMOD-cut-le1*:

$t \in [r, \text{mod } m] \implies$
 $[r, \text{mod } m] \downarrow \leq t = [r, \text{mod } m, (t - r) \text{ div } m]$
 ⟨proof⟩

lemma *iMODb-cut-le*:

$[r, \text{mod } m, c] \downarrow \leq t =$ (
 if $t < r$ then {}
 else if $t < r + m * c$ then $[r, \text{mod } m, (t - r) \text{ div } m]$
 else $[r, \text{mod } m, c]$)
 ⟨proof⟩

lemma *iMODb-cut-le1*:

$t \in [r, \text{mod } m, c] \implies$
 $[r, \text{mod } m, c] \downarrow \leq t = [r, \text{mod } m, (t - r) \text{ div } m]$
 ⟨proof⟩

lemma *iTILL-cut-less*:

$[..n] \downarrow < t =$ (
 if $n < t$ then $[..n]$ else
 if $t = 0$ then {}
 else $[..t - \text{Suc } 0]$)
 ⟨proof⟩

lemma *iTILL-cut-less1*:

$\llbracket t \in [..n]; 0 < t \rrbracket \implies [..n] \downarrow < t = [..t - \text{Suc } 0]$
 ⟨proof⟩

lemma *iFROM-cut-less*:

$[n..] \downarrow < t =$ (
 if $t \leq n$ then {}
 else $[n.., t - \text{Suc } n]$)
 ⟨proof⟩

lemma *iFROM-cut-less1*:

$n < t \implies [n..] \downarrow < t = [n.., t - \text{Suc } n]$
 ⟨proof⟩

lemma *iIN-cut-less*:

$[n.., d] \downarrow < t =$ (
 if $t \leq n$ then {} else
 if $t \leq n + d$ then $[n.., t - \text{Suc } n]$
 else $[n.., d]$)
 ⟨proof⟩

lemma *iIN-cut-less1*:

$\llbracket t \in [n..d]; n < t \rrbracket \implies [n..d] \downarrow < t = [n.., t - \text{Suc } n]$
 <proof>

lemma *iMOD-cut-less*:

$[r, \text{mod } m] \downarrow < t = ($
 if $t \leq r$ then $\{\}$
 else $[r, \text{mod } m, (t - \text{Suc } r) \text{ div } m]$
 <proof>

lemma *iMOD-cut-less1*:

$\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies$
 $[r, \text{mod } m] \downarrow < t = [r, \text{mod } m, (t - r) \text{ div } m - \text{Suc } 0]$
 <proof>

lemma *iMODb-cut-less*:

$[r, \text{mod } m, c] \downarrow < t = ($
 if $t \leq r$ then $\{\}$ else
 if $r + m * c < t$ then $[r, \text{mod } m, c]$
 else $[r, \text{mod } m, (t - \text{Suc } r) \text{ div } m]$
 <proof>

lemma *iMODb-cut-less1*: $\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies$

$[r, \text{mod } m, c] \downarrow < t = [r, \text{mod } m, (t - r) \text{ div } m - \text{Suc } 0]$
 <proof>

lemmas *iT-cut-le* =

iTILL-cut-le
iFROM-cut-le
iIN-cut-le
iMOD-cut-le
iMODb-cut-le

lemmas *iT-cut-le1* =

iTILL-cut-le1
iFROM-cut-le1
iIN-cut-le1
iMOD-cut-le1
iMODb-cut-le1

lemmas *iT-cut-less* =

iTILL-cut-less
iFROM-cut-less
iIN-cut-less
iMOD-cut-less
iMODb-cut-less

lemmas *iT-cut-less1* =
iTILL-cut-less1
iFROM-cut-less1
iIN-cut-less1
iMOD-cut-less1
iMODb-cut-less1

lemmas *iT-cut-le-less* =
iTILL-cut-le
iTILL-cut-less
iFROM-cut-le
iFROM-cut-less
iIN-cut-le
iIN-cut-less
iMOD-cut-le
iMOD-cut-less
iMODb-cut-le
iMODb-cut-less

lemmas *iT-cut-le-less1* =
iTILL-cut-le1
iTILL-cut-less1
iFROM-cut-le1
iFROM-cut-less1
iIN-cut-le1
iIN-cut-less1
iMOD-cut-le1
iMOD-cut-less1
iMODb-cut-le1
iMODb-cut-less1

lemma *iTILL-cut-ge*:
 $[\dots n] \downarrow_{\geq} t = (\text{if } n < t \text{ then } \{\} \text{ else } [t \dots, n-t])$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-ge1*: $t \in [\dots n] \implies [\dots n] \downarrow_{\geq} t = [t \dots, n-t]$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-ge2*: $t \notin [\dots n] \implies [\dots n] \downarrow_{\geq} t = \{\}$
 $\langle \text{proof} \rangle$

lemma *iTILL-cut-greater*:
 $[\dots n] \downarrow_{>} t = (\text{if } n \leq t \text{ then } \{\} \text{ else } [Suc t \dots, n - Suc t])$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-greater1*:
 $t \in [\dots n] \implies t < n \implies [\dots n] \downarrow_{>} t = [Suc t \dots, n - Suc t]$
 $\langle \text{proof} \rangle$

corollary *iTILL-cut-greater2*: $t \notin [\dots n] \implies [\dots n] \downarrow > t = \{\}$
 ⟨proof⟩

lemma *iFROM-cut-ge*:

$[\dots] \downarrow \geq t = (\text{if } n \leq t \text{ then } [t\dots] \text{ else } [n\dots])$

⟨proof⟩

corollary *iFROM-cut-ge1*: $t \in [n\dots] \implies [n\dots] \downarrow \geq t = [t\dots]$

⟨proof⟩

lemma *iFROM-cut-greater*:

$[\dots] \downarrow > t = (\text{if } n \leq t \text{ then } [\text{Suc } t\dots] \text{ else } [n\dots])$

⟨proof⟩

corollary *iFROM-cut-greater1*:

$t \in [n\dots] \implies [n\dots] \downarrow > t = [\text{Suc } t\dots]$

⟨proof⟩

lemma *iIN-cut-ge*:

$[\dots, d] \downarrow \geq t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \text{if } t \leq n+d \text{ then } [t\dots, n+d-t] \text{ else } \{\})$

⟨proof⟩

corollary *iIN-cut-ge1*: $t \in [n\dots, d] \implies$

$[\dots, d] \downarrow \geq t = [t\dots, n+d-t]$

⟨proof⟩

corollary *iIN-cut-ge2*: $t \notin [n\dots, d] \implies$

$[\dots, d] \downarrow \geq t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \{\})$

⟨proof⟩

lemma *iIN-cut-greater*:

$[\dots, d] \downarrow > t = (\text{if } t < n \text{ then } [n\dots, d] \text{ else } \text{if } t < n+d \text{ then } [\text{Suc } t\dots, n+d - \text{Suc } t] \text{ else } \{\})$

⟨proof⟩

corollary *iIN-cut-greater1*:

$\llbracket t \in [n\dots, d]; t < n+d \rrbracket \implies$

$[\dots, d] \downarrow > t = [\text{Suc } t\dots, n+d - \text{Suc } t]$

⟨proof⟩

lemma *mod-cut-greater-aux-t-less*:

$\llbracket 0 < (m::nat); r \leq t \rrbracket \implies$
 $t < t + m - (t - r) \text{ mod } m$
 <proof>

lemma *mod-cut-greater-aux-le-x*:
 $\llbracket (r::nat) \leq t; t < x; x \text{ mod } m = r \text{ mod } m \rrbracket \implies$
 $t + m - (t - r) \text{ mod } m \leq x$
 <proof>

lemma *iMOD-cut-greater*:
 $[r, \text{mod } m] \downarrow > t =$
 if $t < r$ then $[r, \text{mod } m]$ else
 if $m = 0$ then $\{\}$
 else $[t + m - (t - r) \text{ mod } m, \text{mod } m]$
 <proof>

lemma *iMOD-cut-greater1*:
 $t \in [r, \text{mod } m] \implies$
 $[r, \text{mod } m] \downarrow > t =$
 if $m = 0$ then $\{\}$
 else $[t + m, \text{mod } m]$
 <proof>

lemma *iMODb-cut-greater-aux*:
 $\llbracket 0 < m; t < r + m * c; r \leq t \rrbracket \implies$
 $(r + m * c - (t + m - (t - r) \text{ mod } m)) \text{ div } m =$
 $c - \text{Suc } ((t - r) \text{ div } m)$
 <proof>

lemma *iMODb-cut-greater*:
 $[r, \text{mod } m, c] \downarrow > t =$
 if $t < r$ then $[r, \text{mod } m, c]$ else
 if $r + m * c \leq t$ then $\{\}$
 else $[t + m - (t - r) \text{ mod } m, \text{mod } m, c - \text{Suc } ((t - r) \text{ div } m)]$
 <proof>

lemma *iMODb-cut-greater1*:
 $t \in [r, \text{mod } m, c] \implies$
 $[r, \text{mod } m, c] \downarrow > t =$
 if $r + m * c \leq t$ then $\{\}$
 else $[t + m, \text{mod } m, c - \text{Suc } ((t - r) \text{ div } m)]$
 <proof>

lemma *iMOD-cut-ge*:
 $[r, \text{mod } m] \downarrow \geq t =$
 if $t \leq r$ then $[r, \text{mod } m]$ else
 if $m = 0$ then $\{\}$

else $[t + m - \text{Suc} ((t - \text{Suc } r) \bmod m), \bmod m]$
 ⟨proof⟩

lemma *iMOD-cut-ge1*:

$t \in [r, \bmod m] \implies$
 $[r, \bmod m] \downarrow_{\geq} t = [t, \bmod m]$
 ⟨proof⟩

lemma *iMODb-cut-ge*:

$[r, \bmod m, c] \downarrow_{\geq} t = ($
 if $t \leq r$ then $[r, \bmod m, c]$ else
 if $r + m * c < t$ then $\{\}$
 else $[t + m - \text{Suc} ((t - \text{Suc } r) \bmod m), \bmod m, c - (t + m - \text{Suc } r) \text{ div } m]$
 ⟨proof⟩

lemma *iMODb-cut-ge1*:

$t \in [r, \bmod m, c] \implies$
 $[r, \bmod m, c] \downarrow_{\geq} t = ($
 if $r + m * c < t$ then $\{\}$
 else $[t, \bmod m, c - (t - r) \text{ div } m]$
 ⟨proof⟩

lemma *iMOD-0-cut-greater*:

$t \in [r, \bmod 0] \implies [r, \bmod 0] \downarrow_{>} t = \{\}$
 ⟨proof⟩

lemma *iMODb-0-cut-greater*: $t \in [r, \bmod 0, c] \implies$

$[r, \bmod 0, c] \downarrow_{>} t = \{\}$
 ⟨proof⟩

lemmas *iT-cut-ge* =

iTILL-cut-ge
iFROM-cut-ge
iIN-cut-ge
iMOD-cut-ge
iMODb-cut-ge

lemmas *iT-cut-ge1* =

iTILL-cut-ge1
iFROM-cut-ge1
iIN-cut-ge1
iMOD-cut-ge1
iMODb-cut-ge1

lemmas *iT-cut-greater* =

iTILL-cut-greater
iFROM-cut-greater
iIN-cut-greater

iMOD-cut-greater
iMODb-cut-greater

lemmas *iT-cut-greater1* =
iTILL-cut-greater1
iFROM-cut-greater1
iIN-cut-greater1
iMOD-cut-greater1
iMODb-cut-greater1

lemmas *iT-cut-ge-greater* =
iTILL-cut-ge
iTILL-cut-greater
iFROM-cut-ge
iFROM-cut-greater
iIN-cut-ge
iIN-cut-greater
iMOD-cut-ge
iMOD-cut-greater
iMODb-cut-ge
iMODb-cut-greater

lemmas *iT-cut-ge-greater1* =
iTILL-cut-ge1
iTILL-cut-greater1
iFROM-cut-ge1
iFROM-cut-greater1
iIN-cut-ge1
iIN-cut-greater1
iMOD-cut-ge1
iMOD-cut-greater1
iMODb-cut-ge1
iMODb-cut-greater1

1.6 Cardinality of intervals

lemma *iFROM-card*: $\text{card } [n..] = 0$
<proof>

lemma *iTILL-card*: $\text{card } [..n] = \text{Suc } n$
<proof>

lemma *iIN-card*: $\text{card } [n..,d] = \text{Suc } d$
<proof>

lemma *iMOD-0-card*: $\text{card } [r, \text{mod } 0] = \text{Suc } 0$
<proof>

lemma *iMOD-card*: $0 < m \implies \text{card } [r, \text{mod } m] = 0$
<proof>

lemma *iMOD-card-if*: $\text{card } [r, \text{mod } m] = (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } 0)$
 ⟨proof⟩

lemma *iMODb-mod-0-card*: $\text{card } [r, \text{mod } 0, c] = \text{Suc } 0$
 ⟨proof⟩

lemma *iMODb-card*: $0 < m \implies \text{card } [r, \text{mod } m, c] = \text{Suc } c$
 ⟨proof⟩

lemma *iMODb-card-if*:
 $\text{card } [r, \text{mod } m, c] = (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } \text{Suc } c)$
 ⟨proof⟩

lemmas *iT-card* =
iFROM-card
iTILL-card
iIN-card
iMOD-card-if
iMODb-card-if

Cardinality with *icard*

lemma *iFROM-icard*: $\text{icard } [n..] = \infty$
 ⟨proof⟩

lemma *iTILL-icard*: $\text{icard } [..n] = \text{enat } (\text{Suc } n)$
 ⟨proof⟩

lemma *iIN-icard*: $\text{icard } [n..,d] = \text{enat } (\text{Suc } d)$
 ⟨proof⟩

lemma *iMOD-0-icard*: $\text{icard } [r, \text{mod } 0] = \text{eSuc } 0$
 ⟨proof⟩

lemma *iMOD-icard*: $0 < m \implies \text{icard } [r, \text{mod } m] = \infty$
 ⟨proof⟩

lemma *iMOD-icard-if*: $\text{icard } [r, \text{mod } m] = (\text{if } m = 0 \text{ then } \text{eSuc } 0 \text{ else } \infty)$
 ⟨proof⟩

lemma *iMODb-mod-0-icard*: $\text{icard } [r, \text{mod } 0, c] = \text{eSuc } 0$
 ⟨proof⟩

lemma *iMODb-icard*: $0 < m \implies \text{icard } [r, \text{mod } m, c] = \text{enat } (\text{Suc } c)$
 ⟨proof⟩

lemma *iMODb-icard-if*: $\text{icard } [r, \text{mod } m, c] = \text{enat } (\text{if } m = 0 \text{ then } \text{Suc } 0 \text{ else } \text{Suc } c)$
 ⟨proof⟩

lemmas *iT-icard* =
iFROM-icard
iTILL-icard
iIN-icard
iMOD-icard-if
iMODb-icard-if

1.7 Functions *inext* and *iprev* with intervals

lemma

iFROM-inext: $t \in [n..] \implies \text{inext } t [n..] = \text{Suc } t$ **and**
iTILL-inext: $t < n \implies \text{inext } t [..n] = \text{Suc } t$ **and**
iIN-inext: $\llbracket n \leq t; t < n + d \rrbracket \implies \text{inext } t [n..,d] = \text{Suc } t$
⟨proof⟩

lemma

iFROM-iprev': $t \in [n..] \implies \text{iprev } (\text{Suc } t) [n..] = t$ **and**
iFROM-iprev: $n < t \implies \text{iprev } t [n..] = t - \text{Suc } 0$ **and**
iTILL-iprev: $t \in [..n] \implies \text{iprev } t [..n] = t - \text{Suc } 0$ **and**
iIN-iprev: $\llbracket n < t; t \leq n + d \rrbracket \implies \text{iprev } t [n..,d] = t - \text{Suc } 0$ **and**
iIN-iprev': $\llbracket n \leq t; t < n + d \rrbracket \implies \text{iprev } (\text{Suc } t) [n..,d] = t$
⟨proof⟩

lemma *iMOD-inext*: $t \in [r, \text{mod } m] \implies \text{inext } t [r, \text{mod } m] = t + m$
⟨proof⟩

lemma *iMOD-iprev*: $\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies \text{iprev } t [r, \text{mod } m] = t - m$
⟨proof⟩

lemma *iMOD-iprev'*: $t \in [r, \text{mod } m] \implies \text{iprev } (t + m) [r, \text{mod } m] = t$
⟨proof⟩

lemma *iMODb-inext*:

$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies$
 $\text{inext } t [r, \text{mod } m, c] = t + m$
⟨proof⟩

lemma *iMODb-iprev*:

$\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies$
 $\text{iprev } t [r, \text{mod } m, c] = t - m$
⟨proof⟩

lemma *iMODb-iprev'*:

$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies$
 $\text{iprev } (t + m) [r, \text{mod } m, c] = t$
⟨proof⟩

lemmas *iT-inext* =
iFROM-inext

iTILL-inext
iIN-inext
iMOD-inext
iMODb-inext

lemmas *iT-iprev* =

iFROM-iprev'
iFROM-iprev
iTILL-iprev
iIN-iprev
iIN-iprev'
iMOD-iprev
iMOD-iprev'
iMODb-iprev
iMODb-iprev'

lemma *iFROM-inext-if*:

$\text{inext } t [n..] = (\text{if } t \in [n..] \text{ then } \text{Suc } t \text{ else } t)$
 ⟨proof⟩

lemma *iTILL-inext-if*:

$\text{inext } t [..n] = (\text{if } t < n \text{ then } \text{Suc } t \text{ else } t)$
 ⟨proof⟩

lemma *iIN-inext-if*:

$\text{inext } t [n..d] = (\text{if } n \leq t \wedge t < n + d \text{ then } \text{Suc } t \text{ else } t)$
 ⟨proof⟩

lemma *iMOD-inext-if*:

$\text{inext } t [r, \text{mod } m] = (\text{if } t \in [r, \text{mod } m] \text{ then } t + m \text{ else } t)$
 ⟨proof⟩

lemma *iMODb-inext-if*:

$\text{inext } t [r, \text{mod } m, c] =$
 $(\text{if } t \in [r, \text{mod } m, c] \wedge t < r + m * c \text{ then } t + m \text{ else } t)$
 ⟨proof⟩

lemmas *iT-inext-if* =

iFROM-inext-if
iTILL-inext-if
iIN-inext-if
iMOD-inext-if
iMODb-inext-if

lemma *iFROM-iprev-if*:

$\text{iprev } t [n..] = (\text{if } n < t \text{ then } t - \text{Suc } 0 \text{ else } t)$
 ⟨proof⟩

lemma *iTILL-iprev-if*:

$\text{iprev } t [..n] = (\text{if } t \in [..n] \text{ then } t - \text{Suc } 0 \text{ else } t)$
 ⟨proof⟩

lemma *iIN-iprev-if*:

$$\text{iprev } t [n \dots, d] = (\text{if } n < t \wedge t \leq n + d \text{ then } t - \text{Suc } 0 \text{ else } t)$$

<proof>

lemma *iMOD-iprev-if*:

$$\text{iprev } t [r, \text{mod } m] =$$

$$(\text{if } t \in [r, \text{mod } m] \wedge r < t \text{ then } t - m \text{ else } t)$$

<proof>

lemma *iMODb-iprev-if*:

$$\text{iprev } t [r, \text{mod } m, c] =$$

$$(\text{if } t \in [r, \text{mod } m, c] \wedge r < t \text{ then } t - m \text{ else } t)$$

<proof>

lemmas *iT-iprev-if* =

iFROM-iprev-if
iTILL-iprev-if
iIN-iprev-if
iMOD-iprev-if
iMODb-iprev-if

The difference between an element and the next/previous element is constant if the element is different from Min/Max of the interval

lemma *iFROM-inext-diff-const*:

$$t \in [n \dots] \implies \text{inext } t [n \dots] - t = \text{Suc } 0$$

<proof>

lemma *iFROM-iprev-diff-const*:

$$n < t \implies t - \text{iprev } t [n \dots] = \text{Suc } 0$$

<proof>

lemma *iFROM-iprev-diff-const'*:

$$t \in [n \dots] \implies \text{Suc } t - \text{iprev } (\text{Suc } t) [n \dots] = \text{Suc } 0$$

<proof>

lemma *iTILL-inext-diff-const*:

$$t < n \implies \text{inext } t [\dots n] - t = \text{Suc } 0$$

<proof>

lemma *iTILL-iprev-diff-const*:

$$\llbracket t \in [\dots n]; 0 < t \rrbracket \implies t - \text{iprev } t [\dots n] = \text{Suc } 0$$

<proof>

lemma *iIN-inext-diff-const*:

$$\llbracket n \leq t; t < n + d \rrbracket \implies \text{inext } t [n \dots, d] - t = \text{Suc } 0$$

<proof>

lemma *iIN-iprev-diff-const*:

$$\llbracket n < t; t \leq n + d \rrbracket \implies t - \text{iprev } t [n \dots, d] = \text{Suc } 0$$

<proof>

lemma *iIN-iprev-diff-const'*:

$$\llbracket n \leq t; t < n + d \rrbracket \implies \text{Suc } t - \text{iprev } (\text{Suc } t) [n \dots, d] = \text{Suc } 0$$

<proof>

lemma *iMOD-inext-diff-const*:

$$t \in [r, \text{mod } m] \implies \text{inext } t [r, \text{mod } m] - t = m$$

<proof>

lemma *iMOD-iprev-diff-const'*:

$$t \in [r, \text{mod } m] \implies (t + m) - \text{iprev } (t + m) [r, \text{mod } m] = m$$

<proof>

lemma *iMOD-iprev-diff-const*:

$$\llbracket t \in [r, \text{mod } m]; r < t \rrbracket \implies t - \text{iprev } t [r, \text{mod } m] = m$$

<proof>

lemma *iMODb-inext-diff-const*:

$$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies \text{inext } t [r, \text{mod } m, c] - t = m$$

<proof>

lemma *iMODb-iprev-diff-const'*:

$$\llbracket t \in [r, \text{mod } m, c]; t < r + m * c \rrbracket \implies (t + m) - \text{iprev } (t + m) [r, \text{mod } m, c] = m$$

<proof>

lemma *iMODb-iprev-diff-const*:

$$\llbracket t \in [r, \text{mod } m, c]; r < t \rrbracket \implies t - \text{iprev } t [r, \text{mod } m, c] = m$$

<proof>

lemmas *iT-inext-diff-const =*

iFROM-inext-diff-const

iTILL-inext-diff-const

iIN-inext-diff-const

iMOD-inext-diff-const

iMODb-inext-diff-const

lemmas *iT-iprev-diff-const =*

iFROM-iprev-diff-const

iFROM-iprev-diff-const'

iTILL-iprev-diff-const

iIN-iprev-diff-const

iIN-iprev-diff-const'

iMOD-iprev-diff-const'

iMOD-iprev-diff-const

iMODb-iprev-diff-const'

iMODb-iprev-diff-const

1.7.1 Mirroring of intervals

lemma

iIN-mirror-lem: *mirror-lem* $x [n..d] = n + d - x$ **and**

iTILL-mirror-lem: *mirror-lem* $x [..n] = n - x$ **and**

iMODb-mirror-elem: $\text{mirror-elem } x [r, \text{mod } m, c] = r + r + m * c - x$
 ⟨proof⟩

lemma *iMODb-imirror-bounds*:

$r' + m' * c' \leq l + r \implies$
 $\text{imirror-bounds } [r', \text{mod } m', c'] \text{ } l \text{ } r = [l + r - r' - m' * c', \text{mod } m', c']$
 ⟨proof⟩

lemma *iIN-imirror-bounds*:

$n + d \leq l + r \implies \text{imirror-bounds } [n \dots, d] \text{ } l \text{ } r = [l + r - n - d \dots, d]$
 ⟨proof⟩

lemma *iTILL-imirror-bounds*:

$n \leq l + r \implies \text{imirror-bounds } [\dots, n] \text{ } l \text{ } r = [l + r - n \dots, n]$
 ⟨proof⟩

lemmas *iT-imirror-bounds =*

iTILL-imirror-bounds
iIN-imirror-bounds
iMODb-imirror-bounds

lemma *iMODb-imirror-ident*: $\text{imirror } [r, \text{mod } m, c] = [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iIN-imirror-ident*: $\text{imirror } [n \dots, d] = [n \dots, d]$
 ⟨proof⟩

lemma *iTILL-imirror-ident*: $\text{imirror } [\dots, n] = [\dots, n]$
 ⟨proof⟩

lemmas *iT-imirror-ident =*

iTILL-imirror-ident
iIN-imirror-ident
iMODb-imirror-ident

1.7.2 Functions *inext-nth* and *iprev-nth* on intervals

lemma *iFROM-inext-nth* : $[n \dots] \rightarrow a = n + a$
 ⟨proof⟩

lemma *iIN-inext-nth* : $a \leq d \implies [n \dots, d] \rightarrow a = n + a$
 ⟨proof⟩

lemma *iIN-iprev-nth*: $a \leq d \implies [n \dots, d] \leftarrow a = n + d - a$
 ⟨proof⟩

lemma *iIN-inext-nth-if* :

$[n \dots, d] \rightarrow a = (\text{if } a \leq d \text{ then } n + a \text{ else } n + d)$
 ⟨proof⟩

lemma *iIN-iprev-nth-if*:

$$[n\dots,d] \leftarrow a = (\text{if } a \leq d \text{ then } n + d - a \text{ else } n)$$

<proof>

lemma *iTILL-inext-nth* : $a \leq n \implies [\dots n] \rightarrow a = a$

<proof>

lemma *iTILL-inext-nth-if* :

$$[\dots n] \rightarrow a = (\text{if } a \leq n \text{ then } a \text{ else } n)$$

<proof>

lemma *iTILL-iprev-nth*: $a \leq n \implies [\dots n] \leftarrow a = n - a$

<proof>

lemma *iTILL-iprev-nth-if*:

$$[\dots n] \leftarrow a = (\text{if } a \leq n \text{ then } n - a \text{ else } 0)$$

<proof>

lemma *iMOD-inext-nth*: $[r, \text{mod } m] \rightarrow a = r + m * a$

<proof>

lemma *iMODb-inext-nth*: $a \leq c \implies [r, \text{mod } m, c] \rightarrow a = r + m * a$

<proof>

lemma *iMODb-inext-nth-if*:

$$[r, \text{mod } m, c] \rightarrow a = (\text{if } a \leq c \text{ then } r + m * a \text{ else } r + m * c)$$

<proof>

lemma *iMODb-iprev-nth*:

$$a \leq c \implies [r, \text{mod } m, c] \leftarrow a = r + m * (c - a)$$

<proof>

lemma *iMODb-iprev-nth-if*:

$$[r, \text{mod } m, c] \leftarrow a = (\text{if } a \leq c \text{ then } r + m * (c - a) \text{ else } r)$$

<proof>

lemma *iIN-iFROM-inext-nth*:

$$a \leq d \implies [n\dots,d] \rightarrow a = [n\dots] \rightarrow a$$

<proof>

lemma *iIN-iFROM-inext*:

$$a < n + d \implies \text{inext } a [n\dots,d] = \text{inext } a [n\dots]$$

<proof>

lemma *iMOD-iMODb-inext-nth*:

$$a \leq c \implies [r, \text{mod } m, c] \rightarrow a = [r, \text{mod } m] \rightarrow a$$

<proof>

lemma *iMOD-iMODb-inext*:

$$a < r + m * c \implies \text{iNext } a [r, \text{mod } m, c] = \text{iNext } a [r, \text{mod } m]$$

<proof>

lemma *iMOD-inext-nth-Suc-diff*:

$$([r, \text{mod } m] \rightarrow (\text{Suc } n)) - ([r, \text{mod } m] \rightarrow n) = m$$

<proof>

lemma *iMOD-inext-nth-diff*:

$$([r, \text{mod } m] \rightarrow a) - ([r, \text{mod } m] \rightarrow b) = (a - b) * m$$

<proof>

lemma *iMODb-inext-nth-diff*: $\llbracket a \leq c; b \leq c \rrbracket \implies$

$$([r, \text{mod } m, c] \rightarrow a) - ([r, \text{mod } m, c] \rightarrow b) = (a - b) * m$$

<proof>

1.8 Induction with intervals

lemma *iFROM-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [n..]; P k \rrbracket \implies P (\text{Suc } k); a \in [n..] \rrbracket \implies P a$$

<proof>

lemma *iIN-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [n..,d]; k \neq n + d; P k \rrbracket \implies P (\text{Suc } k); a \in [n..,d] \rrbracket \implies P a$$

<proof>

lemma *iTILL-induct*:

$$\llbracket P 0; \bigwedge k. \llbracket k \in [..n]; k \neq n; P k \rrbracket \implies P (\text{Suc } k); a \in [..n] \rrbracket \implies P a$$

<proof>

lemma *iMOD-induct*:

$$\llbracket P r; \bigwedge k. \llbracket k \in [r, \text{mod } m]; P k \rrbracket \implies P (k + m); a \in [r, \text{mod } m] \rrbracket \implies P a$$

<proof>

lemma *iMODb-induct*:

$$\llbracket P r; \bigwedge k. \llbracket k \in [r, \text{mod } m, c]; k \neq r + m * c; P k \rrbracket \implies P (k + m); a \in [r, \text{mod } m, c] \rrbracket \implies P a$$

<proof>

lemma *iIN-rev-induct*:

$$\llbracket P (n + d); \bigwedge k. \llbracket k \in [n..,d]; k \neq n; P k \rrbracket \implies P (k - \text{Suc } 0); a \in [n..,d] \rrbracket \implies P a$$

<proof>

lemma *iTILL-rev-induct*:

$$\llbracket P n; \bigwedge k. \llbracket k \in [..n]; 0 < k; P k \rrbracket \implies P (k - \text{Suc } 0); a \in [..n] \rrbracket \implies P a$$

<proof>

lemma *iMODb-rev-induct*:

$\llbracket P (r + m * c); \bigwedge k. \llbracket k \in [r, \text{mod } m, c]; k \neq r; P k \rrbracket \implies P (k - m); a \in [r, \text{mod } m, c] \rrbracket \implies P a$
 $\langle \text{proof} \rangle$

end

2 Arithmetic operators on natural intervals

theory *IL-IntervalOperators*

imports *IL-Interval*

begin

2.1 Arithmetic operations with intervals

2.1.1 Addition of and multiplication by constants

definition *iT-Plus* :: $iT \Rightarrow Time \Rightarrow iT$ (**infixl** \oplus 55)

where $I \oplus k \equiv (\lambda n. (n + k)) \text{ ' } I$

definition *iT-Mult* :: $iT \Rightarrow Time \Rightarrow iT$ (**infixl** \otimes 55)

where $iT\text{-Mult-def} : I \otimes k \equiv (\lambda n. (n * k)) \text{ ' } I$

lemma *iT-Plus-image-conv*: $I \oplus k = (\lambda n. (n + k)) \text{ ' } I$

$\langle \text{proof} \rangle$

lemma *iT-Mult-image-conv*: $I \otimes k = (\lambda n. (n * k)) \text{ ' } I$

$\langle \text{proof} \rangle$

lemma *iT-Plus-empty*: $\{\} \oplus k = \{\}$

$\langle \text{proof} \rangle$

lemma *iT-Mult-empty*: $\{\} \otimes k = \{\}$

$\langle \text{proof} \rangle$

lemma *iT-Plus-not-empty*: $I \neq \{\} \implies I \oplus k \neq \{\}$

$\langle \text{proof} \rangle$

lemma *iT-Mult-not-empty*: $I \neq \{\} \implies I \otimes k \neq \{\}$

$\langle \text{proof} \rangle$

lemma *iT-Plus-empty-iff*: $(I \oplus k = \{\}) = (I = \{\})$

$\langle \text{proof} \rangle$

lemma *iT-Mult-empty-iff*: $(I \otimes k = \{\}) = (I = \{\})$

$\langle \text{proof} \rangle$

lemma *iT-Plus-mono*: $A \subseteq B \implies A \oplus k \subseteq B \oplus k$
 ⟨proof⟩

lemma *iT-Mult-mono*: $A \subseteq B \implies A \otimes k \subseteq B \otimes k$
 ⟨proof⟩

lemma *iT-Mult-0*: $I \neq \{\} \implies I \otimes 0 = [\dots 0]$
 ⟨proof⟩

corollary *iT-Mult-0-if*: $I \otimes 0 = (\text{if } I = \{\} \text{ then } \{\} \text{ else } [\dots 0])$
 ⟨proof⟩

lemma *iT-Plus-mem-iff*: $x \in (I \oplus k) = (k \leq x \wedge (x - k) \in I)$
 ⟨proof⟩

lemma *iT-Plus-mem-iff2*: $x + k \in (I \oplus k) = (x \in I)$
 ⟨proof⟩

lemma *iT-Mult-mem-iff-0*: $x \in (I \otimes 0) = (I \neq \{\} \wedge x = 0)$
 ⟨proof⟩

lemma *iT-Mult-mem-iff*:
 $0 < k \implies x \in (I \otimes k) = (x \text{ mod } k = 0 \wedge x \text{ div } k \in I)$
 ⟨proof⟩

lemma *iT-Mult-mem-iff2*: $0 < k \implies x * k \in (I \otimes k) = (x \in I)$
 ⟨proof⟩

lemma *iT-Plus-singleton*: $\{a\} \oplus k = \{a + k\}$
 ⟨proof⟩

lemma *iT-Mult-singleton*: $\{a\} \otimes k = \{a * k\}$
 ⟨proof⟩

lemma *iT-Plus-Un*: $(A \cup B) \oplus k = (A \oplus k) \cup (B \oplus k)$
 ⟨proof⟩

lemma *iT-Mult-Un*: $(A \cup B) \otimes k = (A \otimes k) \cup (B \otimes k)$
 ⟨proof⟩

lemma *iT-Plus-Int*: $(A \cap B) \oplus k = (A \oplus k) \cap (B \oplus k)$
 ⟨proof⟩

lemma *iT-Mult-Int*: $0 < k \implies (A \cap B) \otimes k = (A \otimes k) \cap (B \otimes k)$
 ⟨proof⟩

lemma *iT-Plus-image*: $f \text{ ' } I \oplus n = (\lambda x. f x + n) \text{ ' } I$
 ⟨proof⟩

lemma *iT-Mult-image*: $f \text{ ' } I \otimes n = (\lambda x. f x * n) \text{ ' } I$
 ⟨proof⟩

lemma *iT-Plus-commute*: $I \oplus a \oplus b = I \oplus b \oplus a$
 ⟨proof⟩

lemma *iT-Mult-commute*: $I \otimes a \otimes b = I \otimes b \otimes a$
 ⟨proof⟩

lemma *iT-Plus-assoc*: $I \oplus a \oplus b = I \oplus (a + b)$
 ⟨proof⟩

lemma *iT-Mult-assoc*: $I \otimes a \otimes b = I \otimes (a * b)$
 ⟨proof⟩

lemma *iT-Plus-Mult-distrib*: $I \oplus n \otimes m = I \otimes m \oplus n * m$
 ⟨proof⟩⟨proof⟩⟨proof⟩⟨proof⟩⟨proof⟩⟨proof⟩

lemma *iT-Plus-finite-iff*: $\text{finite } (I \oplus k) = \text{finite } I$
 ⟨proof⟩

lemma *iT-Mult-0-finite*: $\text{finite } (I \otimes 0)$
 ⟨proof⟩

lemma *iT-Mult-finite-iff*: $0 < k \implies \text{finite } (I \otimes k) = \text{finite } I$
 ⟨proof⟩

lemma *iT-Plus-Min*: $I \neq \{\} \implies iMin (I \oplus k) = iMin I + k$
 ⟨proof⟩

lemma *iT-Mult-Min*: $I \neq \{\} \implies iMin (I \otimes k) = iMin I * k$
 ⟨proof⟩

lemma *iT-Plus-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max (I \oplus k) = Max I + k$
 ⟨proof⟩

lemma *iT-Mult-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max (I \otimes k) = Max I * k$
 ⟨proof⟩

corollary

iMOD-mult-0: $[r, \text{mod } m] \otimes 0 = [\dots 0]$ **and**

iMODb-mult-0: $[r, \text{mod } m, c] \otimes 0 = [\dots 0]$ **and**

iFROM-mult-0: $[n \dots] \otimes 0 = [\dots 0]$ **and**

iIN-mult-0: $[n \dots, d] \otimes 0 = [\dots 0]$ **and**

iTILL-mult-0: $[\dots n] \otimes 0 = [\dots 0]$

⟨proof⟩

lemmas *iT-mult-0* =

iTILL-mult-0
iFROM-mult-0
iIN-mult-0
iMOD-mult-0
iMODb-mult-0

lemma *iT-Plus-0*: $I \oplus 0 = I$

<proof>

lemma *iT-Mult-1*: $I \otimes \text{Suc } 0 = I$

<proof>

corollary

iFROM-add-Min: $iMin ([n..] \oplus k) = n + k$ **and**
iIN-add-Min: $iMin ([n..,d] \oplus k) = n + k$ **and**
iTILL-add-Min: $iMin ([..n] \oplus k) = k$ **and**
iMOD-add-Min: $iMin ([r, \text{mod } m] \oplus k) = r + k$ **and**
iMODb-add-Min: $iMin ([r, \text{mod } m, c] \oplus k) = r + k$
<proof>

corollary

iFROM-mult-Min: $iMin ([n..] \otimes k) = n * k$ **and**
iIN-mult-Min: $iMin ([n..,d] \otimes k) = n * k$ **and**
iTILL-mult-Min: $iMin ([..n] \otimes k) = 0$ **and**
iMOD-mult-Min: $iMin ([r, \text{mod } m] \otimes k) = r * k$ **and**
iMODb-mult-Min: $iMin ([r, \text{mod } m, c] \otimes k) = r * k$
<proof>

lemmas *iT-add-Min* =

iIN-add-Min
iTILL-add-Min
iFROM-add-Min
iMOD-add-Min
iMODb-add-Min

lemmas *iT-mult-Min* =

iIN-mult-Min
iTILL-mult-Min
iFROM-mult-Min
iMOD-mult-Min
iMODb-mult-Min

lemma *iFROM-add*: $[n..] \oplus k = [n+k..]$

<proof>

lemma *iIN-add*: $[n..,d] \oplus k = [n+k..,d]$

<proof>

lemma *iTILL-add*: $[\dots i] \oplus k = [k \dots, i]$
<proof>

lemma *iMOD-add*: $[r, \text{mod } m] \oplus k = [r + k, \text{mod } m]$
<proof>

lemma *iMODb-add*: $[r, \text{mod } m, c] \oplus k = [r + k, \text{mod } m, c]$
<proof>

lemmas *iT-add* =
iMOD-add
iMODb-add
iFROM-add
iIN-add
iTILL-add
iT-Plus-singleton

lemma *iFROM-mult*: $[n \dots] \otimes k = [n * k, \text{mod } k]$
<proof>

lemma *iIN-mult*: $[n \dots, d] \otimes k = [n * k, \text{mod } k, d]$
<proof>

lemma *iTILL-mult*: $[\dots n] \otimes k = [0, \text{mod } k, n]$
<proof>

lemma *iMOD-mult*: $[r, \text{mod } m] \otimes k = [r * k, \text{mod } m * k]$
<proof>

lemma *iMODb-mult*:
 $[r, \text{mod } m, c] \otimes k = [r * k, \text{mod } m * k, c]$
<proof>

lemmas *iT-mult* =
iTILL-mult
iFROM-mult
iIN-mult
iMOD-mult
iMODb-mult
iT-Mult-singleton

2.1.2 Some conversions between intervals using constant addition and multiplication

lemma *iFROM-conv*: $[n \dots] = UNIV \oplus n$
<proof>

lemma *iIN-conv*: $[n..d] = [\dots d] \oplus n$
 ⟨proof⟩

lemma *iMOD-conv*: $[r, \text{mod } m] = [0..] \otimes m \oplus r$
 ⟨proof⟩

lemma *iMODb-conv*: $[r, \text{mod } m, c] = [\dots c] \otimes m \oplus r$
 ⟨proof⟩

Some examples showing the utility of `iMODb_conv`

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
 ⟨proof⟩

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
 ⟨proof⟩

lemma $[12, \text{mod } 10, 4] = \{12, 22, 32, 42, 52\}$
 ⟨proof⟩

lemma $[r, \text{mod } m, 4] = \{r, r+m, r+2*m, r+3*m, r+4*m\}$
 ⟨proof⟩

lemma $[2, \text{mod } 10, 4] = \{2, 12, 22, 32, 42\}$
 ⟨proof⟩

2.1.3 Subtraction of constants

definition *iT-Plus-neg* :: $iT \Rightarrow \text{Time} \Rightarrow iT$ (**infixl** $\oplus -$ 55) **where**
 $I \oplus - k \equiv \{x. x + k \in I\}$

lemma *iT-Plus-neg-mem-iff*: $(x \in I \oplus - k) = (x + k \in I)$
 ⟨proof⟩

lemma *iT-Plus-neg-mem-iff2*: $k \leq x \implies (x - k \in I \oplus - k) = (x \in I)$
 ⟨proof⟩

lemma *iT-Plus-neg-image-conv*: $I \oplus - k = (\lambda n. (n - k)) \text{ ' } (I \downarrow \geq k)$
 ⟨proof⟩

lemma *iT-Plus-neg-cut-eq*: $t \leq k \implies (I \downarrow \geq t) \oplus - k = I \oplus - k$
 ⟨proof⟩

lemma *iT-Plus-neg-mono*: $A \subseteq B \implies A \oplus - k \subseteq B \oplus - k$
 ⟨proof⟩

lemma *iT-Plus-neg-empty*: $\{\} \oplus - k = \{\}$
 ⟨proof⟩

lemma *iT-Plus-neg-Max-less-empty*:

$\llbracket \text{finite } I; \text{Max } I < k \rrbracket \implies I \oplus - k = \{\}$
 ⟨proof⟩

lemma *iT-Plus-neg-not-empty-iff*: $(I \oplus - k \neq \{\}) = (\exists x \in I. k \leq x)$
 ⟨proof⟩

lemma *iT-Plus-neg-empty-iff*:
 $(I \oplus - k = \{\}) = (I = \{\}) \vee (\text{finite } I \wedge \text{Max } I < k)$
 ⟨proof⟩

lemma *iT-Plus-neg-assoc*: $(I \oplus - a) \oplus - b = I \oplus - (a + b)$
 ⟨proof⟩

lemma *iT-Plus-neg-commute*: $I \oplus - a \oplus - b = I \oplus - b \oplus - a$
 ⟨proof⟩

lemma *iT-Plus-neg-0*: $I \oplus - 0 = I$
 ⟨proof⟩

lemma *iT-Plus-Plus-neg-assoc*: $b \leq a \implies I \oplus a \oplus - b = I \oplus (a - b)$
 ⟨proof⟩

lemma *iT-Plus-Plus-neg-assoc2*: $a \leq b \implies I \oplus a \oplus - b = I \oplus - (b - a)$
 ⟨proof⟩

lemma *iT-Plus-neg-Plus-le-cut-eq*:
 $a \leq b \implies (I \oplus - a) \oplus b = (I \downarrow \geq a) \oplus (b - a)$
 ⟨proof⟩

corollary *iT-Plus-neg-Plus-le-Min-eq*:
 $\llbracket a \leq b; a \leq iMin I \rrbracket \implies (I \oplus - a) \oplus b = I \oplus (b - a)$
 ⟨proof⟩

lemma *iT-Plus-neg-Plus-ge-cut-eq*:
 $b \leq a \implies (I \oplus - a) \oplus b = (I \downarrow \geq a) \oplus - (a - b)$
 ⟨proof⟩

corollary *iT-Plus-neg-Plus-ge-Min-eq*:
 $\llbracket b \leq a; a \leq iMin I \rrbracket \implies (I \oplus - a) \oplus b = I \oplus - (a - b)$
 ⟨proof⟩

lemma *iT-Plus-neg-Mult-distrib*:
 $0 < m \implies I \oplus - n \otimes m = I \otimes m \oplus - n * m$
 ⟨proof⟩

lemma *iT-Plus-neg-Plus-le-inverse*: $k \leq iMin I \implies I \oplus - k \oplus k = I$
 ⟨proof⟩

lemma *iT-Plus-neg-Plus-inverse*: $I \oplus - k \oplus k = I \downarrow \geq k$

<proof>

lemma *iT-Plus-Plus-neg-inverse*: $I \oplus k \oplus - k = I$
<proof>

lemma *iT-Plus-neg-Un*: $(A \cup B) \oplus - k = (A \oplus - k) \cup (B \oplus - k)$
<proof>

lemma *iT-Plus-neg-Int*: $(A \cap B) \oplus - k = (A \oplus - k) \cap (B \oplus - k)$
<proof>

lemma *iT-Plus-neg-Max-singleton*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies I \oplus - \text{Max } I = \{0\}$
<proof>

lemma *iT-Plus-neg-singleton*: $\{a\} \oplus - k = (\text{if } k \leq a \text{ then } \{a - k\} \text{ else } \{\})$
<proof>

corollary *iT-Plus-neg-singleton1*: $k \leq a \implies \{a\} \oplus - k = \{a - k\}$
<proof>

corollary *iT-Plus-neg-singleton2*: $a < k \implies \{a\} \oplus - k = \{\}$
<proof>

lemma *iT-Plus-neg-finite-iff*: $\text{finite } (I \oplus - k) = \text{finite } I$
<proof>

lemma *iT-Plus-neg-Min*:
 $I \oplus - k \neq \{\} \implies i\text{Min } (I \oplus - k) = i\text{Min } (I \downarrow \geq k) - k$
<proof>

lemma *iT-Plus-neg-Max*:
 $\llbracket \text{finite } I; I \oplus - k \neq \{\} \rrbracket \implies \text{Max } (I \oplus - k) = \text{Max } I - k$
<proof>

Subtractions of constants from intervals

lemma *iFROM-add-neg*: $[n \dots] \oplus - k = [n - k \dots]$
<proof>

lemma *iTILL-add-neg*: $[\dots n] \oplus - k = (\text{if } k \leq n \text{ then } [\dots n - k] \text{ else } \{\})$
<proof>

lemma *iTILL-add-neg1*: $k \leq n \implies [\dots n] \oplus - k = [\dots n - k]$
<proof>

lemma *iTILL-add-neg2*: $n < k \implies [\dots n] \oplus - k = \{\}$
<proof>

lemma *iIN-add-neg*:
 $[n \dots, d] \oplus - k = (\text{if } k \leq n \text{ then } [n - k \dots, d])$

else if $k \leq n + d$ then $[\dots n + d - k]$ else $\{\}$
 ⟨proof⟩

lemma *iIN-add-neg1*: $k \leq n \implies [n\dots,d] \oplus -k = [n - k\dots,d]$
 ⟨proof⟩

lemma *iIN-add-neg2*: $\llbracket n \leq k; k \leq n + d \rrbracket \implies [n\dots,d] \oplus -k = [\dots n + d - k]$
 ⟨proof⟩

lemma *iIN-add-neg3*: $n + d < k \implies [n\dots,d] \oplus -k = \{\}$
 ⟨proof⟩

lemma *iMOD-0-add-neg*: $[r, \text{mod } 0] \oplus -k = \{r\} \oplus -k$
 ⟨proof⟩

lemma *iMOD-gr0-add-neg*:

$0 < m \implies$
 $[r, \text{mod } m] \oplus -k = ($
 if $k \leq r$ then $[r - k, \text{mod } m]$
 else $[(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m]$
 ⟨proof⟩

lemma *iMOD-add-neg*:

$[r, \text{mod } m] \oplus -k = ($
 if $k \leq r$ then $[r - k, \text{mod } m]$
 else if $0 < m$ then $[(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m, \text{mod } m]$ else $\{\}$
 ⟨proof⟩

corollary *iMOD-add-neg1*:

$k \leq r \implies [r, \text{mod } m] \oplus -k = [r - k, \text{mod } m]$
 ⟨proof⟩

lemma *iMOD-add-neg2*:

$\llbracket 0 < m; r < k \rrbracket \implies [r, \text{mod } m] \oplus -k = [(m + r \text{ mod } m - k \text{ mod } m) \text{ mod } m,$
 $\text{mod } m]$
 ⟨proof⟩

lemma *iMODb-mod-0-add-neg*: $[r, \text{mod } 0, c] \oplus -k = \{r\} \oplus -k$
 ⟨proof⟩

lemma *iMODb-add-neg*:

$[r, \text{mod } m, c] \oplus -k = ($
 if $k \leq r$ then $[r - k, \text{mod } m, c]$
 else

$$\text{if } k \leq r + m * c \text{ then}$$

$$[(m + r \bmod m - k \bmod m) \bmod m, \bmod m, (r + m * c - k) \text{ div } m]$$

$$\text{else } \{\}$$
 (proof)

lemma *iMODb-add-neg'*:

$$[r, \bmod m, c] \oplus - k = ($$

$$\text{if } k \leq r \text{ then } [r - k, \bmod m, c]$$

$$\text{else if } k \leq r + m * c \text{ then}$$

$$\text{if } k \bmod m \leq r \bmod m$$

$$\text{then } [r \bmod m - k \bmod m, \bmod m, c + r \text{ div } m - k \text{ div } m]$$

$$\text{else } [m + r \bmod m - k \bmod m, \bmod m, c + r \text{ div } m - \text{Suc } (k \text{ div } m)]$$

$$\text{else } \{\}$$
 (proof)

corollary *iMODb-add-neg1*:

$$k \leq r \implies [r, \bmod m, c] \oplus - k = [r - k, \bmod m, c]$$
 (proof)

corollary *iMODb-add-neg2*:

$$[[r < k; k \leq r + m * c] \implies$$

$$[r, \bmod m, c] \oplus - k =$$

$$[(m + r \bmod m - k \bmod m) \bmod m, \bmod m, (r + m * c - k) \text{ div } m]$$
 (proof)

corollary *iMODb-add-neg2-mod-le*:

$$[[r < k; k \leq r + m * c; k \bmod m \leq r \bmod m] \implies$$

$$[r, \bmod m, c] \oplus - k =$$

$$[r \bmod m - k \bmod m, \bmod m, c + r \text{ div } m - k \text{ div } m]$$
 (proof)

corollary *iMODb-add-neg2-mod-less*:

$$[[r < k; k \leq r + m * c; r \bmod m < k \bmod m] \implies$$

$$[r, \bmod m, c] \oplus - k =$$

$$[m + r \bmod m - k \bmod m, \bmod m, c + r \text{ div } m - \text{Suc } (k \text{ div } m)]$$
 (proof)

lemma *iMODb-add-neg3*: $r + m * c < k \implies [r, \bmod m, c] \oplus - k = \{\}$
 (proof)

lemmas *iT-add-neg* =

iFROM-add-neg
iIN-add-neg
iTILL-add-neg
iMOD-add-neg
iMODb-add-neg
iT-Plus-neg-singleton

2.1.4 Subtraction of intervals from constants

definition *iT-Minus* :: *Time* \Rightarrow *iT* \Rightarrow *iT* (**infixl** \ominus 55)

where $k \ominus I \equiv \{x. x \leq k \wedge (k - x) \in I\}$

lemma *iT-Minus-mem-iff*: $(x \in k \ominus I) = (x \leq k \wedge k - x \in I)$

<proof>

lemma *iT-Minus-mono*: $A \subseteq B \Longrightarrow k \ominus A \subseteq k \ominus B$

<proof>

lemma *iT-Minus-image-conv*: $k \ominus I = (\lambda x. k - x) ` (I \downarrow \leq k)$

<proof>

lemma *iT-Minus-cut-eq*: $k \leq t \Longrightarrow k \ominus (I \downarrow \leq t) = k \ominus I$

<proof>

lemma *iT-Minus-Minus-cut-eq*: $k \ominus (k \ominus (I \downarrow \leq k)) = I \downarrow \leq k$

<proof>

lemma $10 \ominus [\dots 3] = [7\dots, 3]$

<proof>

lemma *iT-Minus-empty*: $k \ominus \{\} = \{\}$

<proof>

lemma *iT-Minus-0*: $k \ominus \{0\} = \{k\}$

<proof>

lemma *iT-Minus-bound*: $x \in k \ominus I \Longrightarrow x \leq k$

<proof>

lemma *iT-Minus-finite*: *finite* $(k \ominus I)$

<proof>

lemma *iT-Minus-less-Min-empty*: $k < iMin I \Longrightarrow k \ominus I = \{\}$

<proof>

lemma *iT-Minus-Min-singleton*: $I \neq \{\} \Longrightarrow (iMin I) \ominus I = \{0\}$

<proof>

lemma *iT-Minus-empty-iff*: $(k \ominus I = \{\}) = (I = \{\} \vee k < iMin I)$

<proof>

lemma *iT-Minus-imirror-conv*:

$k \ominus I = imirror (I \downarrow \leq k) \oplus k \oplus - (iMin I + Max (I \downarrow \leq k))$

<proof>

lemma *iT-Minus-imirror-conv'*:

$$k \ominus I = \text{imirror } (I \downarrow \leq k) \oplus k \oplus - (iMin (I \downarrow \leq k) + Max (I \downarrow \leq k))$$

<proof>

lemma *iT-Minus-Max*:

$$\llbracket I \neq \{\} ; iMin I \leq k \rrbracket \implies Max (k \ominus I) = k - (iMin I)$$

<proof>

lemma *iT-Minus-Min*:

$$\llbracket I \neq \{\} ; iMin I \leq k \rrbracket \implies iMin (k \ominus I) = k - (Max (I \downarrow \leq k))$$

<proof>

lemma *iT-Minus-Minus-eq*: $\llbracket \text{finite } I ; Max I \leq k \rrbracket \implies k \ominus (k \ominus I) = I$

<proof>

lemma *iT-Minus-Minus-eq2*: $I \subseteq [\dots k] \implies k \ominus (k \ominus I) = I$

<proof>

lemma *iT-Minus-Minus*: $a \ominus (b \ominus I) = (I \downarrow \leq b) \oplus a \oplus - b$

<proof>

lemma *iT-Minus-Plus-empty*: $k < n \implies k \ominus (I \oplus n) = \{\}$

<proof>

lemma *iT-Minus-Plus-commute*: $n \leq k \implies k \ominus (I \oplus n) = (k - n) \ominus I$

<proof>

lemma *iT-Minus-Plus-cut-assoc*: $(k \ominus I) \oplus n = (k + n) \ominus (I \downarrow \leq k)$

<proof>

lemma *iT-Minus-Plus-assoc*:

$$\llbracket \text{finite } I ; Max I \leq k \rrbracket \implies (k \ominus I) \oplus n = (k + n) \ominus I$$

<proof>

lemma *iT-Minus-Plus-assoc2*:

$$I \subseteq [\dots k] \implies (k \ominus I) \oplus n = (k + n) \ominus I$$

<proof>

lemma *iT-Minus-Un*: $k \ominus (A \cup B) = (k \ominus A) \cup (k \ominus B)$

<proof>

lemma *iT-Minus-Int*: $k \ominus (A \cap B) = (k \ominus A) \cap (k \ominus B)$

<proof>

lemma *iT-Minus-singleton*: $k \ominus \{a\} = (\text{if } a \leq k \text{ then } \{k - a\} \text{ else } \{\})$

<proof>

corollary *iT-Minus-singleton1*: $a \leq k \implies k \ominus \{a\} = \{k - a\}$

<proof>

corollary *iT-Minus-singleton2*: $k < a \implies k \ominus \{a\} = \{\}$

<proof>

lemma *iMOD-sub*:

$k \ominus [r, \text{mod } m] =$
 (if $r \leq k$ then $[(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$ else $\{\}$)
<proof>

corollary *iMOD-sub1*:

$r \leq k \implies k \ominus [r, \text{mod } m] = [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$
<proof>

corollary *iMOD-sub2*: $k < r \implies k \ominus [r, \text{mod } m] = \{\}$
<proof>

lemma *iTILL-sub*: $k \ominus [\dots n] = (\text{if } n \leq k \text{ then } [k - n \dots, n] \text{ else } [\dots k])$
<proof>

corollary *iTILL-sub1*: $n \leq k \implies k \ominus [\dots n] = [k - n \dots, n]$
<proof>

corollary *iTILL-sub2*: $k \leq n \implies k \ominus [\dots n] = [\dots k]$
<proof>

lemma *iMODb-sub*:

$k \ominus [r, \text{mod } m, c] =$
 (if $r + m * c \leq k$ then $[k - r - m * c, \text{mod } m, c]$ else
 if $r \leq k$ then $[(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$ else $\{\}$)
<proof>

corollary *iMODb-sub1*:

$\llbracket r \leq k; k \leq r + m * c \rrbracket \implies$
 $k \ominus [r, \text{mod } m, c] = [(k - r) \text{ mod } m, \text{mod } m, (k - r) \text{ div } m]$
<proof>

corollary *iMODb-sub2*: $k < r \implies k \ominus [r, \text{mod } m, c] = \{\}$
<proof>

corollary *iMODb-sub3*:

$r + m * c \leq k \implies k \ominus [r, \text{mod } m, c] = [k - r - m * c, \text{mod } m, c]$
<proof>

lemma *iFROM-sub*: $k \ominus [n \dots] = (\text{if } n \leq k \text{ then } [\dots k - n] \text{ else } \{\})$
<proof>

corollary *iFROM-sub1*: $n \leq k \implies k \ominus [n \dots] = [\dots k - n]$
<proof>

corollary *iFROM-sub-empty*: $k < n \implies k \ominus [n..] = \{\}$
 ⟨proof⟩

lemma *iIN-sub*:
 $k \ominus [n..,d] =$
 if $n + d \leq k$ then $[k - (n + d)..,d]$
 else if $n \leq k$ then $[..k - n]$ else $\{\}$
 ⟨proof⟩

lemma *iIN-sub1*: $n + d \leq k \implies k \ominus [n..,d] = [k - (n + d)..,d]$
 ⟨proof⟩

lemma *iIN-sub2*: $\llbracket n \leq k; k \leq n + d \rrbracket \implies k \ominus [n..,d] = [..k - n]$
 ⟨proof⟩

lemma *iIN-sub3*: $k < n \implies k \ominus [n..,d] = \{\}$
 ⟨proof⟩

lemmas *iT-sub* =
iFROM-sub
iIN-sub
iTILL-sub
iMOD-sub
iMODb-sub
iT-Minus-singleton

2.1.5 Division of intervals by constants

Monotonicity and injectivity of arithmetic operators

lemma *iMOD-div-right-strict-mono-on*:
 $\llbracket 0 < k; k \leq m \rrbracket \implies \text{strict-mono-on } (\lambda x. x \text{ div } k) [r, \text{ mod } m]$
 ⟨proof⟩

corollary *iMOD-div-right-inj-on*:
 $\llbracket 0 < k; k \leq m \rrbracket \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{ mod } m]$
 ⟨proof⟩

lemma *iMOD-mult-div-right-inj-on*:
 $\text{inj-on } (\lambda x. x \text{ div } (k::\text{nat})) [r, \text{ mod } (k * m)]$
 ⟨proof⟩

lemma *iMOD-mult-div-right-inj-on2*:
 $m \text{ mod } k = 0 \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{ mod } m]$
 ⟨proof⟩

lemma *iMODb-div-right-strict-mono-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{strict-mono-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$
 ⟨proof⟩

corollary *iMODb-div-right-inj-on*:

$\llbracket 0 < k; k \leq m \rrbracket \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$
 ⟨proof⟩

lemma *iMODb-mult-div-right-inj-on*:

$\text{inj-on } (\lambda x. x \text{ div } (k::\text{nat})) [r, \text{mod } (k * m), c]$
 ⟨proof⟩

corollary *iMODb-mult-div-right-inj-on2*:

$m \text{ mod } k = 0 \implies \text{inj-on } (\lambda x. x \text{ div } k) [r, \text{mod } m, c]$
 ⟨proof⟩

definition *iT-Div* :: $iT \Rightarrow Time \Rightarrow iT$ (**infixl** \circ 55)

where $I \circ k \equiv (\lambda n. (n \text{ div } k)) \text{ ' } I$

lemma *iT-Div-image-conv*: $I \circ k = (\lambda n. (n \text{ div } k)) \text{ ' } I$
 ⟨proof⟩

lemma *iT-Div-mono*: $A \subseteq B \implies A \circ k \subseteq B \circ k$
 ⟨proof⟩

lemma *iT-Div-empty*: $\{\} \circ k = \{\}$

⟨proof⟩

lemma *iT-Div-not-empty*: $I \neq \{\} \implies I \circ k \neq \{\}$

⟨proof⟩

lemma *iT-Div-empty-iff*: $(I \circ k = \{\}) = (I = \{\})$

⟨proof⟩

lemma *iT-Div-0*: $I \neq \{\} \implies I \circ 0 = [\dots 0]$

⟨proof⟩

corollary *iT-Div-0-if*: $I \circ 0 = (\text{if } I = \{\} \text{ then } \{\} \text{ else } [\dots 0])$

⟨proof⟩

corollary

iFROM-div-0: $[n\dots] \circ 0 = [\dots 0]$ **and**

iTILL-div-0: $[\dots n] \circ 0 = [\dots 0]$ **and**

iIN-div-0: $[n\dots, d] \circ 0 = [\dots 0]$ **and**

iMOD-div-0: $[r, \text{mod } m] \circ 0 = [\dots 0]$ **and**

iMODb-div-0: $[r, \text{mod } m, c] \circ 0 = [\dots 0]$

⟨proof⟩

lemmas *iT-div-0* =

iTILL-div-0

iFROM-div-0

iIN-div-0
iMOD-div-0
iMODb-div-0

lemma *iT-Div-1*: $I \circ \text{Suc } 0 = I$
 ⟨proof⟩

lemma *iT-Div-mem-iff-0*: $x \in (I \circ 0) = (I \neq \{\}) \wedge x = 0$
 ⟨proof⟩

lemma *iT-Div-mem-iff*:
 $0 < k \implies x \in (I \circ k) = (\exists y \in I. y \text{ div } k = x)$
 ⟨proof⟩

lemma *iT-Div-mem-iff2*:
 $0 < k \implies x \text{ div } k \in (I \circ k) = (\exists y \in I. y \text{ div } k = x \text{ div } k)$
 ⟨proof⟩

lemma *iT-Div-mem-iff-Int*:
 $0 < k \implies x \in (I \circ k) = (I \cap [x * k \dots k - \text{Suc } 0] \neq \{\})$
 ⟨proof⟩

lemma *iT-Div-imp-mem*:
 $0 < k \implies x \in I \implies x \text{ div } k \in (I \circ k)$
 ⟨proof⟩

lemma *iT-Div-singleton*: $\{a\} \circ k = \{a \text{ div } k\}$
 ⟨proof⟩

lemma *iT-Div-Un*: $(A \cup B) \circ k = (A \circ k) \cup (B \circ k)$
 ⟨proof⟩

lemma *iT-Div-insert*: $(\text{insert } n \ I) \circ k = \text{insert } (n \text{ div } k) \ (I \circ k)$
 ⟨proof⟩

lemma *not-iT-Div-Int*: $\neg (\forall k \ A \ B. (A \cap B) \circ k = (A \circ k) \cap (B \circ k))$
 ⟨proof⟩

lemma *subset-iT-Div-Int*: $A \subseteq B \implies (A \cap B) \circ k = (A \circ k) \cap (B \circ k)$
 ⟨proof⟩

lemma *iFROM-iT-Div-Int*:
 $\llbracket 0 < k; n \leq iMin \ A \rrbracket \implies (A \cap [n \dots]) \circ k = (A \circ k) \cap ([n \dots] \circ k)$
 ⟨proof⟩

lemma *iIN-iT-Div-Int*:
 $\llbracket 0 < k; n \leq iMin \ A; \forall x \in A. x \text{ div } k \leq (n + d) \text{ div } k \implies x \leq n + d \rrbracket \implies$

$$(A \cap [n..d]) \otimes k = (A \otimes k) \cap ([n..d] \otimes k)$$

<proof>

corollary *iTILL-iT-Div-Int:*

$$\llbracket 0 < k; \forall x \in A. x \text{ div } k \leq n \text{ div } k \longrightarrow x \leq n \rrbracket \implies$$

$$(A \cap [..n]) \otimes k = (A \otimes k) \cap ([..n] \otimes k)$$

<proof>

lemma *iIN-iT-Div-Int-mod-0:*

$$\llbracket 0 < k; n \bmod k = 0; \forall x \in A. x \text{ div } k \leq (n + d) \text{ div } k \longrightarrow x \leq n + d \rrbracket \implies$$

$$(A \cap [n..d]) \otimes k = (A \otimes k) \cap ([n..d] \otimes k)$$

<proof>

lemma *mod-partition-iT-Div-Int:*

$$\llbracket 0 < k; 0 < d \rrbracket \implies$$

$$(A \cap [n * k..d * k - \text{Suc } 0]) \otimes k =$$

$$(A \otimes k) \cap ([n * k..d * k - \text{Suc } 0] \otimes k)$$

<proof><proof>

corollary *mod-partition-iT-Div-Int2:*

$$\llbracket 0 < k; 0 < d; n \bmod k = 0; d \bmod k = 0 \rrbracket \implies$$

$$(A \cap [n..d - \text{Suc } 0]) \otimes k =$$

$$(A \otimes k) \cap ([n..d - \text{Suc } 0] \otimes k)$$

<proof>

corollary *mod-partition-iT-Div-Int-one-segment:*

$$0 < k \implies$$

$$(A \cap [n * k..k - \text{Suc } 0]) \otimes k = (A \otimes k) \cap ([n * k..k - \text{Suc } 0] \otimes k)$$

<proof>

corollary *mod-partition-iT-Div-Int-one-segment2:*

$$\llbracket 0 < k; n \bmod k = 0 \rrbracket \implies$$

$$(A \cap [n..k - \text{Suc } 0]) \otimes k = (A \otimes k) \cap ([n..k - \text{Suc } 0] \otimes k)$$

<proof>

lemma *iT-Div-assoc:* $I \otimes a \otimes b = I \otimes (a * b)$

<proof>

lemma *iT-Div-commute:* $I \otimes a \otimes b = I \otimes b \otimes a$

<proof>

lemma *iT-Mult-Div-self:* $0 < k \implies I \otimes k \otimes k = I$

<proof>

lemma *iT-Mult-Div:*

$$\llbracket 0 < d; k \bmod d = 0 \rrbracket \implies I \otimes k \otimes d = I \otimes (k \text{ div } d)$$

<proof>

lemma *iT-Div-Mult-self:*

$$0 < k \implies I \otimes k \otimes k = \{y. \exists x \in I. y = x - x \bmod k\}$$

<proof>

lemma *iT-Plus-Div-distrib-mod-less*:

$$\forall x \in I. x \bmod m + n \bmod m < m \implies I \oplus n \otimes m = I \otimes m \oplus n \operatorname{div} m$$

<proof>

corollary *iT-Plus-Div-distrib-mod-0*:

$$n \bmod m = 0 \implies I \oplus n \otimes m = I \otimes m \oplus n \operatorname{div} m$$

<proof>

lemma *iT-Div-Min*: $I \neq \{\}$ $\implies iMin (I \otimes k) = iMin I \operatorname{div} k$
<proof>

corollary

iFROM-div-Min: $iMin ([n..] \otimes k) = n \operatorname{div} k$ **and**
iIN-div-Min: $iMin ([n..,d] \otimes k) = n \operatorname{div} k$ **and**
iTILL-div-Min: $iMin ([..n] \otimes k) = 0$ **and**
iMOD-div-Min: $iMin ([r, \operatorname{mod} m] \otimes k) = r \operatorname{div} k$ **and**
iMODb-div-Min: $iMin ([r, \operatorname{mod} m, c] \otimes k) = r \operatorname{div} k$

<proof>

lemmas *iT-div-Min* =

iFROM-div-Min
iIN-div-Min
iTILL-div-Min
iMOD-div-Min
iMODb-div-Min

lemma *iT-Div-Max*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies Max (I \otimes k) = Max I \operatorname{div} k$
<proof>

corollary

iIN-div-Max: $Max ([n..,d] \otimes k) = (n + d) \operatorname{div} k$ **and**
iTILL-div-Max: $Max ([..n] \otimes k) = n \operatorname{div} k$ **and**
iMODb-div-Max: $Max ([r, \operatorname{mod} m, c] \otimes k) = (r + m * c) \operatorname{div} k$

<proof>

lemma *iT-Div-0-finite*: $\text{finite} (I \otimes 0)$
<proof>

lemma *iT-Div-infinite-iff*: $0 < k \implies \text{infinite} (I \otimes k) = \text{infinite } I$
<proof>

lemma *iT-Div-finite-iff*: $0 < k \implies \text{finite} (I \otimes k) = \text{finite } I$
<proof>

lemma *iFROM-div*: $0 < k \implies [n..] \otimes k = [n \operatorname{div} k..]$
<proof>

lemma *iIN-div*:

$$0 < k \implies [n..,d] \otimes k = [n \operatorname{div} k.., d \operatorname{div} k + (n \bmod k + d \bmod k) \operatorname{div} k]$$

$\langle \text{proof} \rangle$

corollary *iIN-div-if*:

$$0 < k \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, d \operatorname{div} k + (\text{if } n \operatorname{mod} k + d \operatorname{mod} k < k \text{ then } 0 \text{ else } \operatorname{Suc} 0)]$$

$\langle \text{proof} \rangle$

corollary *iIN-div-eq1*:

$$\llbracket 0 < k; n \operatorname{mod} k + d \operatorname{mod} k < k \rrbracket \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, d \operatorname{div} k]$$

$\langle \text{proof} \rangle$

corollary *iIN-div-eq2*:

$$\llbracket 0 < k; k \leq n \operatorname{mod} k + d \operatorname{mod} k \rrbracket \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, \operatorname{Suc} (d \operatorname{div} k)]$$

$\langle \text{proof} \rangle$

corollary *iIN-div-mod-eq-0*:

$$\llbracket 0 < k; n \operatorname{mod} k = 0 \rrbracket \implies [n \dots, d] \odot k = [n \operatorname{div} k \dots, d \operatorname{div} k]$$

$\langle \text{proof} \rangle$

lemma *iTILL-div*:

$$0 < k \implies [\dots n] \odot k = [\dots n \operatorname{div} k]$$

$\langle \text{proof} \rangle$

lemma *iMOD-div-ge*:

$$\llbracket 0 < m; m \leq k \rrbracket \implies [r, \operatorname{mod} m] \odot k = [r \operatorname{div} k \dots]$$

$\langle \text{proof} \rangle$

corollary *iMOD-div-self*:

$$0 < m \implies [r, \operatorname{mod} m] \odot m = [r \operatorname{div} m \dots]$$

$\langle \text{proof} \rangle$

lemma *iMOD-div*:

$$\llbracket 0 < k; m \operatorname{mod} k = 0 \rrbracket \implies [r, \operatorname{mod} m] \odot k = [r \operatorname{div} k, \operatorname{mod} (m \operatorname{div} k)]$$

$\langle \text{proof} \rangle$

lemma *iMODb-div-self*:

$$0 < m \implies [r, \operatorname{mod} m, c] \odot m = [r \operatorname{div} m \dots, c]$$

$\langle \text{proof} \rangle$

lemma *iMODb-div-ge*:

$$\llbracket 0 < m; m \leq k \rrbracket \implies [r, \operatorname{mod} m, c] \odot k = [r \operatorname{div} k \dots, (r + m * c) \operatorname{div} k - r \operatorname{div} k]$$

$\langle \text{proof} \rangle$

corollary *iMODb-div-ge-if*:

$$\llbracket 0 < m; m \leq k \rrbracket \implies$$

$[r, \text{mod } m, c] \odot k =$
 $[r \text{ div } k \dots, m * c \text{ div } k + (\text{if } r \text{ mod } k + m * c \text{ mod } k < k \text{ then } 0 \text{ else } \text{Suc } 0)]$
 <proof>

lemma *iMODb-div*:

$\llbracket 0 < k; m \text{ mod } k = 0 \rrbracket \implies$
 $[r, \text{mod } m, c] \odot k = [r \text{ div } k, \text{mod } (m \text{ div } k), c]$
 <proof>

lemmas *iT-div =*

iTILL-div
iFROM-div
iIN-div
iMOD-div
iMODb-div
iT-Div-singleton

This lemma is valid for all $k \leq m$, i. e., also for k with $m \text{ mod } k \neq (0::'a)$.

lemma *iMODb-div-unique*:

$\llbracket 0 < k; k \leq m; k \leq c; [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k \rrbracket \implies$
 $r' = r \text{ div } k \wedge m' = m \text{ div } k \wedge c' = c$
 <proof>

lemma *iMODb-div-mod-gr0-is-0-not-ex0*:

$\llbracket 0 < k; k < m; 0 < m \text{ mod } k; k \leq c; r \text{ mod } k = 0 \rrbracket \implies$
 $\neg(\exists r' m' c'. [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k)$
 <proof>

lemma *iMODb-div-mod-gr0-not-ex--arith-aux1*:

$\llbracket (0::\text{nat}) < k; k < m; 0 < x1 \rrbracket \implies$
 $x1 * m + x2 - x \text{ mod } k + x3 + x \text{ mod } k = x1 * m + x2 + x3$
 <proof>

lemma *iMODb-div-mod-gr0-not-ex*:

$\llbracket 0 < k; k < m; 0 < m \text{ mod } k; k \leq c \rrbracket \implies$
 $\neg(\exists r' m' c'. [r', \text{mod } m', c'] = [r, \text{mod } m, c] \odot k)$
 <proof>

lemma *iMOD-div-eq-imp-iMODb-div-eq*:

$\llbracket 0 < k; k \leq m; [r', \text{mod } m'] = [r, \text{mod } m] \odot k \rrbracket \implies$
 $[r', \text{mod } m', c] = [r, \text{mod } m, c] \odot k$
 <proof>

lemma *iMOD-div-unique*:

$\llbracket 0 < k; k \leq m; [r', \text{mod } m'] = [r, \text{mod } m] \odot k \rrbracket \implies$
 $r' = r \text{ div } k \wedge m' = m \text{ div } k$

$\langle \text{proof} \rangle$

lemma *iMOD-div-mod-gr0-not-ex*:

$$\llbracket 0 < k; k < m; 0 < m \bmod k \rrbracket \implies \\ \neg (\exists r' m'. [r', \bmod m'] = [r, \bmod m] \otimes k)$$

$\langle \text{proof} \rangle$

2.2 Interval cut operators with arithmetic interval operators

lemma

$$\begin{aligned} \textit{iT-Plus-cut-le2}: \quad & (I \oplus k) \downarrow_{\leq} (t + k) = (I \downarrow_{\leq} t) \oplus k \textbf{ and} \\ \textit{iT-Plus-cut-less2}: \quad & (I \oplus k) \downarrow_{<} (t + k) = (I \downarrow_{<} t) \oplus k \textbf{ and} \\ \textit{iT-Plus-cut-ge2}: \quad & (I \oplus k) \downarrow_{\geq} (t + k) = (I \downarrow_{\geq} t) \oplus k \textbf{ and} \\ \textit{iT-Plus-cut-greater2}: \quad & (I \oplus k) \downarrow_{>} (t + k) = (I \downarrow_{>} t) \oplus k \end{aligned}$$

$\langle \text{proof} \rangle$

lemma *iT-Plus-cut-le*:

$$(I \oplus k) \downarrow_{\leq} t = (\textit{if } t < k \textit{ then } \{\} \textit{ else } I \downarrow_{\leq} (t - k) \oplus k)$$

$\langle \text{proof} \rangle$

lemma *iT-Plus-cut-less*: $(I \oplus k) \downarrow_{<} t = I \downarrow_{<} (t - k) \oplus k$

$\langle \text{proof} \rangle$

lemma *iT-Plus-cut-ge*: $(I \oplus k) \downarrow_{\geq} t = I \downarrow_{\geq} (t - k) \oplus k$

$\langle \text{proof} \rangle$

lemma *iT-Plus-cut-greater*:

$$(I \oplus k) \downarrow_{>} t = (\textit{if } t < k \textit{ then } I \oplus k \textit{ else } I \downarrow_{>} (t - k) \oplus k)$$

$\langle \text{proof} \rangle$

lemma

$$\begin{aligned} \textit{iT-Mult-cut-le2}: \quad & 0 < k \implies (I \otimes k) \downarrow_{\leq} (t * k) = (I \downarrow_{\leq} t) \otimes k \textbf{ and} \\ \textit{iT-Mult-cut-less2}: \quad & 0 < k \implies (I \otimes k) \downarrow_{<} (t * k) = (I \downarrow_{<} t) \otimes k \textbf{ and} \\ \textit{iT-Mult-cut-ge2}: \quad & 0 < k \implies (I \otimes k) \downarrow_{\geq} (t * k) = (I \downarrow_{\geq} t) \otimes k \textbf{ and} \\ \textit{iT-Mult-cut-greater2}: \quad & 0 < k \implies (I \otimes k) \downarrow_{>} (t * k) = (I \downarrow_{>} t) \otimes k \end{aligned}$$

$\langle \text{proof} \rangle$

lemma *iT-Mult-cut-le*:

$$0 < k \implies (I \otimes k) \downarrow_{\leq} t = (I \downarrow_{\leq} (t \textit{ div } k)) \otimes k$$

$\langle \text{proof} \rangle$

lemma *iT-Mult-cut-less*:

$$0 < k \implies (I \otimes k) \downarrow_{<} t = \\ (\textit{if } t \bmod k = 0 \textit{ then } (I \downarrow_{<} (t \textit{ div } k)) \textit{ else } I \downarrow_{<} \textit{Suc } (t \textit{ div } k)) \otimes k$$

$\langle \text{proof} \rangle$

lemma *iT-Mult-cut-greater*:

$$0 < k \implies (I \otimes k) \downarrow_{>} t = (I \downarrow_{>} (t \textit{ div } k)) \otimes k$$

<proof>

lemma *iT-Mult-cut-ge:*

$$0 < k \implies (I \otimes k) \downarrow_{\geq} t =$$

$$(if\ t\ mod\ k = 0\ then\ (I \downarrow_{\geq} (t\ div\ k))\ else\ I \downarrow_{\geq} Suc\ (t\ div\ k)) \otimes k$$

<proof>

lemma *iT-Plus-neg-cut-le2:* $k \leq t \implies (I \oplus - k) \downarrow_{\leq} (t - k) = (I \downarrow_{\leq} t) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-less2:* $(I \oplus - k) \downarrow_{<} (t - k) = (I \downarrow_{<} t) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-ge2:* $(I \oplus - k) \downarrow_{\geq} (t - k) = (I \downarrow_{\geq} t) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-greater2:* $k \leq t \implies (I \oplus - k) \downarrow_{>} (t - k) = (I \downarrow_{>} t) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-le:* $(I \oplus - k) \downarrow_{\leq} t = I \downarrow_{\leq} (t + k) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-less:* $(I \oplus - k) \downarrow_{<} t = I \downarrow_{<} (t + k) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-ge:* $(I \oplus - k) \downarrow_{\geq} t = I \downarrow_{\geq} (t + k) \oplus - k$

<proof>

lemma *iT-Plus-neg-cut-greater:* $(I \oplus - k) \downarrow_{>} t = I \downarrow_{>} (t + k) \oplus - k$

<proof>

lemma *iT-Minus-cut-le2:* $t \leq k \implies (k \ominus I) \downarrow_{\leq} (k - t) = k \ominus (I \downarrow_{\geq} t)$

<proof>

lemma *iT-Minus-cut-less2:* $(k \ominus I) \downarrow_{<} (k - t) = k \ominus (I \downarrow_{>} t)$

<proof>

lemma *iT-Minus-cut-ge2:* $(k \ominus I) \downarrow_{\geq} (k - t) = k \ominus (I \downarrow_{\leq} t)$

<proof>

lemma *iT-Minus-cut-greater2:* $t \leq k \implies (k \ominus I) \downarrow_{>} (k - t) = k \ominus (I \downarrow_{<} t)$

<proof>

lemma *iT-Minus-cut-le:* $(k \ominus I) \downarrow_{\leq} t = k \ominus (I \downarrow_{\geq} (k - t))$

<proof>

lemma *iT-Minus-cut-less:*

$(k \ominus I) \downarrow < t = (\text{if } t \leq k \text{ then } k \ominus (I \downarrow > (k - t)) \text{ else } k \ominus I)$
 ⟨proof⟩

lemma *iT-Minus-cut-ge*:

$(k \ominus I) \downarrow \geq t = (\text{if } t \leq k \text{ then } k \ominus (I \downarrow \leq (k - t)) \text{ else } \{\})$
 ⟨proof⟩

lemma *iT-Minus-cut-greater*: $(k \ominus I) \downarrow > t = k \ominus (I \downarrow < (k - t))$
 ⟨proof⟩

lemma *iT-Div-cut-le*:

$0 < k \implies (I \circledast k) \downarrow \leq t = I \downarrow \leq (t * k + (k - \text{Suc } 0)) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-less*:

$0 < k \implies (I \circledast k) \downarrow < t = I \downarrow < (t * k) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-ge*:

$0 < k \implies (I \circledast k) \downarrow \geq t = I \downarrow \geq (t * k) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-greater*:

$0 < k \implies (I \circledast k) \downarrow > t = I \downarrow > (t * k + (k - \text{Suc } 0)) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-le2*:

$0 < k \implies (I \circledast k) \downarrow \leq (t \text{ div } k) = I \downarrow \leq (t - t \text{ mod } k + (k - \text{Suc } 0)) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-less2*:

$0 < k \implies (I \circledast k) \downarrow < (t \text{ div } k) = I \downarrow < (t - t \text{ mod } k) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-ge2*:

$0 < k \implies (I \circledast k) \downarrow \geq (t \text{ div } k) = I \downarrow \geq (t - t \text{ mod } k) \circledast k$
 ⟨proof⟩

lemma *iT-Div-cut-greater2*:

$0 < k \implies (I \circledast k) \downarrow > (t \text{ div } k) = I \downarrow > (t - t \text{ mod } k + (k - \text{Suc } 0)) \circledast k$
 ⟨proof⟩

2.3 *inext* and *iprev* with interval operators

lemma *iT-Plus-inext*: $\text{inext } (n + k) (I \oplus k) = (\text{inext } n I) + k$
 ⟨proof⟩

lemma *iT-Plus-iprev*: $iprev (n + k) (I \oplus k) = (iprev n I) + k$
 ⟨proof⟩

lemma *iT-Plus-inext2*: $k \leq n \implies inext n (I \oplus k) = (inext (n - k) I) + k$
 ⟨proof⟩

lemma *iT-Plus-prev2*: $k \leq n \implies iprev n (I \oplus k) = (iprev (n - k) I) + k$
 ⟨proof⟩

lemma *iT-Mult-inext*: $inext (n * k) (I \otimes k) = (inext n I) * k$
 ⟨proof⟩

lemma *iT-Mult-iprev*: $iprev (n * k) (I \otimes k) = (iprev n I) * k$
 ⟨proof⟩

lemma *iT-Mult-inext2-if*:
 $inext n (I \otimes k) = (if n \text{ mod } k = 0 \text{ then } (inext (n \text{ div } k) I) * k \text{ else } n)$
 ⟨proof⟩

lemma *iT-Mult-iprev2-if*:
 $iprev n (I \otimes k) = (if n \text{ mod } k = 0 \text{ then } (iprev (n \text{ div } k) I) * k \text{ else } n)$
 ⟨proof⟩

corollary *iT-Mult-inext2*:
 $n \text{ mod } k = 0 \implies inext n (I \otimes k) = (inext (n \text{ div } k) I) * k$
 ⟨proof⟩

corollary *iT-Mult-iprev2*:
 $n \text{ mod } k = 0 \implies iprev n (I \otimes k) = (iprev (n \text{ div } k) I) * k$
 ⟨proof⟩

lemma *iT-Plus-neg-inext*:
 $k \leq n \implies inext (n - k) (I \oplus - k) = inext n I - k$
 ⟨proof⟩

lemma *iT-Plus-neg-iprev*:
 $iprev (n - k) (I \oplus - k) = iprev n (I \downarrow \geq k) - k$
 ⟨proof⟩

corollary *iT-Plus-neg-inext2*: $inext n (I \oplus - k) = inext (n + k) I - k$
 ⟨proof⟩

corollary *iT-Plus-neg-iprev2*: $iprev n (I \oplus - k) = iprev (n + k) (I \downarrow \geq k) - k$
 ⟨proof⟩

lemma *iT-Minus-inext*:
 $\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies inext (k - n) (k \ominus I) = k - iprev n I$
 ⟨proof⟩

corollary *iT-Minus-inext2*:

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{inext } n (k \ominus I) = k - \text{iprev } (k - n) I$$

<proof>

lemma *iT-Minus-iprev*:

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{iprev } (k - n) (k \ominus I) = k - \text{inext } n (I \downarrow \leq k)$$

<proof>

lemma *iT-Minus-iprev2*:

$$\llbracket k \ominus I \neq \{\}; n \leq k \rrbracket \implies \text{iprev } n (k \ominus I) = k - \text{inext } (k - n) (I \downarrow \leq k)$$

<proof>

lemma *iT-Plus-inext-nth*: $I \neq \{\} \implies (I \oplus k) \rightarrow n = (I \rightarrow n) + k$

<proof>

lemma *iT-Plus-iprev-nth*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies (I \oplus k) \leftarrow n = (I \leftarrow n) + k$

<proof>

lemma *iT-Mult-inext-nth*: $I \neq \{\} \implies (I \otimes k) \rightarrow n = (I \rightarrow n) * k$

<proof>

lemma *iT-Mult-iprev-nth*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies (I \otimes k) \leftarrow n = (I \leftarrow n) * k$

<proof>

lemma *iT-Plus-neg-inext-nth*:

$$I \oplus - k \neq \{\} \implies (I \oplus - k) \rightarrow n = (I \downarrow \geq k \rightarrow n) - k$$

<proof>

lemma *iT-Plus-neg-iprev-nth*:

$$\llbracket \text{finite } I; I \oplus - k \neq \{\} \rrbracket \implies (I \oplus - k) \leftarrow n = (I \downarrow \geq k \leftarrow n) - k$$

<proof>

lemma *iT-Minus-inext-nth*:

$$k \ominus I \neq \{\} \implies (k \ominus I) \rightarrow n = k - ((I \downarrow \leq k) \leftarrow n)$$

<proof>

lemma *iT-Minus-iprev-nth*:

$$k \ominus I \neq \{\} \implies (k \ominus I) \leftarrow n = k - ((I \downarrow \leq k) \rightarrow n)$$

<proof>

lemma *iT-Div-ge-inext-nth*:

$$\llbracket I \neq \{\}; \forall x \in I. \forall y \in I. x < y \implies x + k \leq y \rrbracket \implies$$

$$(I \oslash k) \rightarrow n = (I \rightarrow n) \text{ div } k$$

<proof>

lemma *iT-Div-mod-inext-nth*:

$$\llbracket I \neq \{\}; \forall x \in I. \forall y \in I. x \text{ mod } k = y \text{ mod } k \rrbracket \implies$$

$(I \otimes k) \rightarrow n = (I \rightarrow n) \text{ div } k$
 ⟨proof⟩

lemma *iT-Div-ge-iprev-nth*:

$\llbracket \text{finite } I; I \neq \{\}; \forall x \in I. \forall y \in I. x < y \longrightarrow x + k \leq y \rrbracket \implies$
 $(I \otimes k) \leftarrow n = (I \leftarrow n) \text{ div } k$
 ⟨proof⟩

lemma *iT-Div-mod-iprev-nth*:

$\llbracket \text{finite } I; I \neq \{\}; \forall x \in I. \forall y \in I. x \bmod k = y \bmod k \rrbracket \implies$
 $(I \otimes k) \leftarrow n = (I \leftarrow n) \text{ div } k$
 ⟨proof⟩

2.4 Cardinality of intervals with interval operators

lemma *iT-Plus-card*: $\text{card } (I \oplus k) = \text{card } I$

⟨proof⟩

lemma *iT-Mult-card*: $0 < k \implies \text{card } (I \otimes k) = \text{card } I$

⟨proof⟩

lemma *iT-Plus-neg-card*: $\text{card } (I \oplus - k) = \text{card } (I \downarrow \geq k)$

⟨proof⟩

lemma *iT-Plus-neg-card-le*: $\text{card } (I \oplus - k) \leq \text{card } I$

⟨proof⟩

lemma *iT-Minus-card*: $\text{card } (k \ominus I) = \text{card } (I \downarrow \leq k)$

⟨proof⟩

lemma *iT-Minus-card-le*: $\text{finite } I \implies \text{card } (k \ominus I) \leq \text{card } I$

⟨proof⟩

lemma *iT-Div-0-card-if*:

$\text{card } (I \otimes 0) = (\text{if } I = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 ⟨proof⟩

lemma *Int-empty-sum*:

$(\sum_{k \leq (n :: \text{nat})} \text{if } \{\} \cap (I k) = \{\} \text{ then } 0 \text{ else } \text{Suc } 0) = 0$
 ⟨proof⟩

lemma *iT-Div-mod-partition-card*:

$\text{card } (I \cap [n * d \dots, d - \text{Suc } 0] \otimes d) =$
 $(\text{if } I \cap [n * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$
 ⟨proof⟩

lemma *iT-Div-conv-count*:

$0 < d \implies I \otimes d = \{k. I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\}$
 ⟨proof⟩

lemma *iT-Div-conv-count2*:

$\llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $I \odot d = \{k. k \leq n \wedge I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\}$
 ⟨proof⟩

lemma *mod-partition-count-Suc*:

$\{k. k \leq \text{Suc } n \wedge I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\} =$
 $\{k. k \leq n \wedge I \cap [k * d \dots, d - \text{Suc } 0] \neq \{\}\} \cup$
 $(\text{if } I \cap [\text{Suc } n * d \dots, d - \text{Suc } 0] \neq \{\} \text{ then } \{\text{Suc } n\} \text{ else } \{\})$
 ⟨proof⟩

lemma *iT-Div-card*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } (I \odot d) = (\sum_{k \leq n} \dots)$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0$
 ⟨proof⟩

corollary *iT-Div-card-Suc*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } (I \odot d) = (\sum_{k < \text{Suc } n} \dots)$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0$
 ⟨proof⟩

corollary *iT-Div-Max-card*: $\llbracket 0 < d; \text{finite } I \rrbracket \implies$

$\text{card } (I \odot d) = (\sum_{k \leq \text{Max } I \text{ div } d} \dots)$
 $\text{if } I \cap [k * d \dots, d - \text{Suc } 0] = \{\} \text{ then } 0 \text{ else } \text{Suc } 0$
 ⟨proof⟩

lemma *iT-Div-card-le*: $0 < k \implies \text{card } (I \odot k) \leq \text{card } I$

⟨proof⟩

lemma *iT-Div-card-inj-on*:

$\text{inj-on } (\lambda n. n \text{ div } k) I \implies \text{card } (I \odot k) = \text{card } I$
 ⟨proof⟩

lemma *mod-Suc'*:

$0 < n \implies \text{Suc } m \text{ mod } n = (\text{if } m \text{ mod } n < n - \text{Suc } 0 \text{ then } \text{Suc } (m \text{ mod } n) \text{ else } 0)$
 ⟨proof⟩

lemma *div-Suc*:

$0 < n \implies \text{Suc } m \text{ div } n = (\text{if } \text{Suc } (m \text{ mod } n) = n \text{ then } \text{Suc } (m \text{ div } n) \text{ else } m \text{ div } n)$
 ⟨proof⟩

lemma *div-Suc'*:

$0 < n \implies \text{Suc } m \text{ div } n = (\text{if } m \text{ mod } n < n - \text{Suc } 0 \text{ then } m \text{ div } n \text{ else } \text{Suc } (m \text{ div } n))$
 ⟨proof⟩

lemma *iT-Div-card-ge-aux*:

$\bigwedge I. \llbracket 0 < d; \text{finite } I; \text{Max } I \text{ div } d \leq n \rrbracket \implies$
 $\text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{card } (I \oslash d)$
 ⟨proof⟩

lemma *iT-Div-card-ge*:

$\text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq \text{card } (I \oslash d)$
 ⟨proof⟩

corollary *iT-Div-card-ge-div*: $\text{card } I \text{ div } d \leq \text{card } (I \oslash d)$

⟨proof⟩

There is no better lower bound function f for $i \oslash d$ with $\text{card } i$ and d as arguments.

lemma *iT-Div-card-ge--is-maximal-lower-bound*:

$\forall I d. \text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0) \leq f (\text{card } I) d \wedge$
 $f (\text{card } I) d \leq \text{card } (I \oslash d) \implies$
 $f (\text{card } (I::\text{nat set})) d = \text{card } I \text{ div } d + (\text{if } \text{card } I \text{ mod } d = 0 \text{ then } 0 \text{ else } \text{Suc } 0)$
 ⟨proof⟩

lemma *iT-Plus-icard*: $\text{icard } (I \oplus k) = \text{icard } I$

⟨proof⟩

lemma *iT-Mult-icard*: $0 < k \implies \text{icard } (I \otimes k) = \text{icard } I$

⟨proof⟩

lemma *iT-Plus-neg-icard*: $\text{icard } (I \oplus - k) = \text{icard } (I \downarrow \geq k)$

⟨proof⟩

lemma *iT-Plus-neg-icard-le*: $\text{icard } (I \oplus - k) \leq \text{icard } I$

⟨proof⟩

lemma *iT-Minus-icard*: $\text{icard } (k \ominus I) = \text{icard } (I \downarrow \leq k)$

⟨proof⟩

lemma *iT-Minus-icard-le*: $\text{icard } (k \ominus I) \leq \text{icard } I$

⟨proof⟩

lemma *iT-Div-0-icard-if*: $\text{icard } (I \oslash 0) = \text{enat } (\text{if } I = \{\} \text{ then } 0 \text{ else } \text{Suc } 0)$

⟨proof⟩

lemma *iT-Div-mod-partition-icard*:

$\text{icard } (I \cap [n * d.., d - \text{Suc } 0] \oslash d) =$

$enat (if I \cap [n * d \dots, d - Suc 0] = \{\} then 0 else Suc 0)$
 <proof>

lemma *iT-Div-icard*:

$\llbracket 0 < d; finite I \implies Max I div d \leq n \rrbracket \implies$
 $icard (I \odot d) =$
 $(if finite I then enat (\sum k \leq n. if I \cap [k * d \dots, d - Suc 0] = \{\} then 0 else Suc$
 $0) else \infty)$
 <proof>

corollary *iT-Div-Max-icard*: $0 < d \implies$

$icard (I \odot d) = (if finite I$
 $then enat (\sum k \leq Max I div d. if I \cap [k * d \dots, d - Suc 0] = \{\} then 0 else Suc$
 $0) else \infty)$
 <proof>

lemma *iT-Div-icard-le*: $0 < k \implies icard (I \odot k) \leq icard I$

<proof>

lemma *iT-Div-icard-inj-on*: $inj-on (\lambda n. n div k) I \implies icard (I \odot k) = icard I$

<proof>

lemma *iT-Div-icard-ge*: $icard I div (enat d) + enat (if icard I mod (enat d) = 0$
 $then 0 else Suc 0) \leq icard (I \odot d)$

<proof>

corollary *iT-Div-icard-ge-div*: $icard I div (enat d) \leq icard (I \odot d)$

<proof>

lemma *iT-Div-icard-ge-is-maximal-lower-bound*:

$\forall I d. icard I div (enat d) + enat (if icard I mod (enat d) = 0 then 0 else Suc 0)$
 $\leq f (icard I) d \wedge$
 $f (icard I) d \leq icard (I \odot d) \implies$
 $f (icard (I :: nat set)) d =$
 $icard I div (enat d) + enat (if icard I mod (enat d) = 0 then 0 else Suc 0)$
 <proof>

2.5 Results about sets of intervals

2.5.1 Set of intervals without and with empty interval

definition *iFROM-UN-set* :: $(nat set) set$

where $iFROM-UN-set \equiv \bigcup n. \{[n \dots]\}$

definition *iTILL-UN-set* :: $(nat set) set$

where $iTILL-UN-set \equiv \bigcup n. \{[...n]\}$

definition *iIN-UN-set* :: $(nat set) set$

where $iIN-UN-set \equiv \bigcup n d. \{[n \dots, d]\}$

definition *iMOD-UN-set* :: (nat set) set
 where *iMOD-UN-set* $\equiv \bigcup r m. \{[r, \text{mod } m]\}$

definition *iMODb-UN-set* :: (nat set) set
 where *iMODb-UN-set* $\equiv \bigcup r m c. \{[r, \text{mod } m, c]\}$

definition *iFROM-set* :: (nat set) set
 where *iFROM-set* $\equiv \{[n..] \mid n. \text{True}\}$

definition *iTILL-set* :: (nat set) set
 where *iTILL-set* $\equiv \{[..n] \mid n. \text{True}\}$

definition *iIN-set* :: (nat set) set
 where *iIN-set* $\equiv \{[n..,d] \mid n d. \text{True}\}$

definition *iMOD-set* :: (nat set) set
 where *iMOD-set* $\equiv \{[r, \text{mod } m] \mid r m. \text{True}\}$

definition *iMODb-set* :: (nat set) set
 where *iMODb-set* $\equiv \{[r, \text{mod } m, c] \mid r m c. \text{True}\}$

definition *iMOD2-set* :: (nat set) set
 where *iMOD2-set* $\equiv \{[r, \text{mod } m] \mid r m. 2 \leq m\}$

definition *iMODb2-set* :: (nat set) set
 where *iMODb2-set* $\equiv \{[r, \text{mod } m, c] \mid r m c. 2 \leq m \wedge 1 \leq c\}$

definition *iMOD2-UN-set* :: (nat set) set
 where *iMOD2-UN-set* $\equiv \bigcup r. \bigcup m \in \{2..\}. \{[r, \text{mod } m]\}$

definition *iMODb2-UN-set* :: (nat set) set
 where *iMODb2-UN-set* $\equiv \bigcup r. \bigcup m \in \{2..\}. \bigcup c \in \{1..\}. \{[r, \text{mod } m, c]\}$

lemmas *i-set-defs* =
iFROM-set-def *iTILL-set-def* *iIN-set-def*
iMOD-set-def *iMODb-set-def*
iMOD2-set-def *iMODb2-set-def*

lemmas *i-UN-set-defs* =
iFROM-UN-set-def *iTILL-UN-set-def* *iIN-UN-set-def*
iMOD-UN-set-def *iMODb-UN-set-def*
iMOD2-UN-set-def *iMODb2-UN-set-def*

lemma *iFROM-set-UN-set-eq*: *iFROM-set* = *iFROM-UN-set*
 ⟨*proof*⟩

lemma

iTILL-set-UN-set-eq: $iTILL\text{-set} = iTILL\text{-UN-set}$ **and**
iIN-set-UN-set-eq: $iIN\text{-set} = iIN\text{-UN-set}$ **and**
iMOD-set-UN-set-eq: $iMOD\text{-set} = iMOD\text{-UN-set}$ **and**
iMODb-set-UN-set-eq: $iMODb\text{-set} = iMODb\text{-UN-set}$

<proof>

lemma *iMOD2-set-UN-set-eq*: $iMOD2\text{-set} = iMOD2\text{-UN-set}$

<proof>

lemma *iMODb2-set-UN-set-eq*: $iMODb2\text{-set} = iMODb2\text{-UN-set}$

<proof>

lemmas *i-set-i-UN-set-sets-eq* =

iFROM-set-UN-set-eq
iTILL-set-UN-set-eq
iIN-set-UN-set-eq
iMOD-set-UN-set-eq
iMODb-set-UN-set-eq
iMOD2-set-UN-set-eq
iMODb2-set-UN-set-eq

lemma *iMOD2-set-iMOD-set-subset*: $iMOD2\text{-set} \subseteq iMOD\text{-set}$

<proof>

lemma *iMODb2-set-iMODb-set-subset*: $iMODb2\text{-set} \subseteq iMODb\text{-set}$

<proof>

definition *i-set* :: (nat set) set

where $i\text{-set} \equiv iFROM\text{-set} \cup iTILL\text{-set} \cup iIN\text{-set} \cup iMOD\text{-set} \cup iMODb\text{-set}$

definition *i-UN-set* :: (nat set) set

where $i\text{-UN-set} \equiv iFROM\text{-UN-set} \cup iTILL\text{-UN-set} \cup iIN\text{-UN-set} \cup iMOD\text{-UN-set} \cup iMODb\text{-UN-set}$

Minimal definitions for *i-set* and *i-set*

definition *i-set-min* :: (nat set) set

where $i\text{-set-min} \equiv iFROM\text{-set} \cup iIN\text{-set} \cup iMOD2\text{-set} \cup iMODb2\text{-set}$

definition *i-UN-set-min* :: (nat set) set

where $i\text{-UN-set-min} \equiv iFROM\text{-UN-set} \cup iIN\text{-UN-set} \cup iMOD2\text{-UN-set} \cup iMODb2\text{-UN-set}$

definition *i-set0* :: (nat set) set

where $i\text{-set0} \equiv \text{insert } \{\} \ i\text{-set}$

lemma *i-set-i-UN-set-eq*: $i\text{-set} = i\text{-UN-set}$

<proof>

lemma *i-set-min-i-UN-set-min-eq*: $i\text{-set-min} = i\text{-UN-set-min}$
 ⟨proof⟩

lemma *i-set-min-eq*: $i\text{-set} = i\text{-set-min}$
 ⟨proof⟩

corollary *i-UN-set-i-UN-min-set-eq*: $i\text{-UN-set} = i\text{-UN-set-min}$
 ⟨proof⟩

lemma *i-set-min-is-minimal-let*:
 let $s1 = i\text{FROM-set}$; $s2 = i\text{IN-set}$; $s3 = i\text{MOD2-set}$; $s4 = i\text{MODb2-set}$ in
 $s1 \cap s2 = \{\}$ \wedge $s1 \cap s3 = \{\}$ \wedge $s1 \cap s4 = \{\}$ \wedge
 $s2 \cap s3 = \{\}$ \wedge $s2 \cap s4 = \{\}$ \wedge $s3 \cap s4 = \{\}$
 ⟨proof⟩

lemmas *i-set-min-is-minimal* = *i-set-min-is-minimal-let*[simplified]

inductive-set *i-set-ind*:: (nat set) set

where

i-set-ind-FROM[intro!]: $[n..] \in i\text{-set-ind}$
 | *i-set-ind-TILL*[intro!]: $[..n] \in i\text{-set-ind}$
 | *i-set-ind-IN*[intro!]: $[n..,d] \in i\text{-set-ind}$
 | *i-set-ind-MOD*[intro!]: $[r, \text{mod } m] \in i\text{-set-ind}$
 | *i-set-ind-MODb*[intro!]: $[r, \text{mod } m, c] \in i\text{-set-ind}$

inductive-set *i-set0-ind* :: (nat set) set

where

i-set0-ind-empty[intro!]: $\{\} \in i\text{-set0-ind}$
 | *i-set0-ind-i-set*[intro!]: $I \in i\text{-set-ind} \implies I \in i\text{-set0-ind}$

The introduction rule *i-set0-ind-i-set* is not declared a safe introduction rule, because it would disturb the correct usage of the *safe* method.

lemma *i-set-ind-subset-i-set0-ind*: $i\text{-set-ind} \subseteq i\text{-set0-ind}$
 ⟨proof⟩

lemma

i-set0-ind-FROM[intro!]: $[n..] \in i\text{-set0-ind}$ **and**
i-set0-ind-TILL[intro!]: $[..n] \in i\text{-set0-ind}$ **and**
i-set0-ind-IN[intro!]: $[n..,d] \in i\text{-set0-ind}$ **and**
i-set0-ind-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set0-ind}$ **and**
i-set0-ind-MODb[intro!]: $[r, \text{mod } m, c] \in i\text{-set0-ind}$
 ⟨proof⟩

lemmas *i-set0-ind-intros2* =

i-set0-ind-empty
i-set0-ind-FROM
i-set0-ind-TILL

i-set0-ind-IN
i-set0-ind-MOD
i-set0-ind-MODb

lemma *i-set-i-set-ind-eq*: $i\text{-set} = i\text{-set-ind}$
 $\langle \text{proof} \rangle$

lemma *i-set0-i-set0-ind-eq*: $i\text{-set0} = i\text{-set0-ind}$
 $\langle \text{proof} \rangle$

lemma *i-set-imp-not-empty*: $I \in i\text{-set} \implies I \neq \{\}$
 $\langle \text{proof} \rangle$

lemma *i-set0-i-set-mem-conv*: $(I \in i\text{-set0}) = (I \in i\text{-set} \vee I = \{\})$
 $\langle \text{proof} \rangle$

lemma *i-set-i-set0-mem-conv*: $(I \in i\text{-set}) = (I \in i\text{-set0} \wedge I \neq \{\})$
 $\langle \text{proof} \rangle$

lemma *i-set0-i-set-conv*: $i\text{-set0} - \{\{\}\} = i\text{-set}$
 $\langle \text{proof} \rangle$

corollary *i-set-subset-i-set0*: $i\text{-set} \subseteq i\text{-set0}$
 $\langle \text{proof} \rangle$

lemma *i-set-singleton*: $\{a\} \in i\text{-set}$
 $\langle \text{proof} \rangle$

lemma *i-set0-singleton*: $\{a\} \in i\text{-set0}$
 $\langle \text{proof} \rangle$

corollary

i-set-FROM[intro!]: $[n..] \in i\text{-set}$ **and**
i-set-TILL[intro!]: $[..n] \in i\text{-set}$ **and**
i-set-IN[intro!]: $[n..,d] \in i\text{-set}$ **and**
i-set-MOD[intro!]: $[r, \text{mod } m] \in i\text{-set}$ **and**
i-set-MODb[intro!]: $[r, \text{mod } m, c] \in i\text{-set}$
 $\langle \text{proof} \rangle$

lemmas *i-set-intros* =

i-set-FROM
i-set-TILL
i-set-IN
i-set-MOD
i-set-MODb

lemma
i-set0-empty[intro!]: $\{\} \in i\text{-set0}$ **and**
i-set0-FROM[intro!]: $[n..] \in i\text{-set0}$ **and**

$i\text{-set0-TILL}[\text{intro!}] : [\dots n] \in i\text{-set0}$ **and**
 $i\text{-set0-IN}[\text{intro!}] : [n \dots d] \in i\text{-set0}$ **and**
 $i\text{-set0-MOD}[\text{intro!}] : [r, \text{mod } m] \in i\text{-set0}$ **and**
 $i\text{-set0-MODb}[\text{intro!}] : [r, \text{mod } m, c] \in i\text{-set0}$
 ⟨proof⟩

lemmas $i\text{-set0-intros} =$

$i\text{-set0-empty}$
 $i\text{-set0-FROM}$
 $i\text{-set0-TILL}$
 $i\text{-set0-IN}$
 $i\text{-set0-MOD}$
 $i\text{-set0-MODb}$

lemma $i\text{-set-infinite-as-iMOD}$:

$\llbracket I \in i\text{-set}; \text{infinite } I \rrbracket \implies \exists r m. I = [r, \text{mod } m]$
 ⟨proof⟩

lemma $i\text{-set-finite-as-iMODb}$:

$\llbracket I \in i\text{-set}; \text{finite } I \rrbracket \implies \exists r m c. I = [r, \text{mod } m, c]$
 ⟨proof⟩

lemma $i\text{-set-as-iMOD-iMODb}$:

$I \in i\text{-set} \implies (\exists r m. I = [r, \text{mod } m]) \vee (\exists r m c. I = [r, \text{mod } m, c])$
 ⟨proof⟩

2.5.2 Interval sets are closed under cutting

lemma $i\text{-set-cut-le-ge-closed-disj}$:

$\llbracket I \in i\text{-set}; t \in I; \text{cut-op} = (\downarrow \leq) \vee \text{cut-op} = (\downarrow \geq) \rrbracket \implies$
 $\text{cut-op } I t \in i\text{-set}$
 ⟨proof⟩

corollary

$i\text{-set-cut-le-closed}$: $\llbracket I \in i\text{-set}; t \in I \rrbracket \implies I \downarrow \leq t \in i\text{-set}$ **and**
 $i\text{-set-cut-ge-closed}$: $\llbracket I \in i\text{-set}; t \in I \rrbracket \implies I \downarrow \geq t \in i\text{-set}$
 ⟨proof⟩

lemmas $i\text{-set-cut-le-ge-closed} = i\text{-set-cut-le-closed } i\text{-set-cut-ge-closed}$

lemma $i\text{-set0-cut-closed-disj}$:

$\llbracket I \in i\text{-set0};$
 $\text{cut-op} = (\downarrow \leq) \vee \text{cut-op} = (\downarrow \geq) \vee$
 $\text{cut-op} = (\downarrow <) \vee \text{cut-op} = (\downarrow >) \rrbracket \implies$
 $\text{cut-op } I t \in i\text{-set0}$
 ⟨proof⟩

corollary

i-set0-cut-le-closed: $I \in i\text{-set0} \implies I \downarrow \leq t \in i\text{-set0}$ **and**
i-set0-cut-less-closed: $I \in i\text{-set0} \implies I \downarrow < t \in i\text{-set0}$ **and**
i-set0-cut-ge-closed: $I \in i\text{-set0} \implies I \downarrow \geq t \in i\text{-set0}$ **and**
i-set0-cut-greater-closed: $I \in i\text{-set0} \implies I \downarrow > t \in i\text{-set0}$
 ⟨proof⟩

lemmas *i-set0-cut-closed* =
i-set0-cut-le-closed
i-set0-cut-less-closed
i-set0-cut-ge-closed
i-set0-cut-greater-closed

2.5.3 Interval sets are closed under addition and multiplication

lemma *i-set-Plus-closed*: $I \in i\text{-set} \implies I \oplus k \in i\text{-set}$
 ⟨proof⟩

lemma *i-set-Mult-closed*: $I \in i\text{-set} \implies I \otimes k \in i\text{-set}$
 ⟨proof⟩

lemma *i-set0-Plus-closed*: $I \in i\text{-set0} \implies I \oplus k \in i\text{-set0}$
 ⟨proof⟩

lemma *i-set0-Mult-closed*: $I \in i\text{-set0} \implies I \otimes k \in i\text{-set0}$
 ⟨proof⟩

2.5.4 Interval sets are closed with certain conditions under subtraction

lemma *i-set-Plus-neg-closed*:
 [$I \in i\text{-set}; \exists x \in I. k \leq x$] $\implies I \oplus -k \in i\text{-set}$
 ⟨proof⟩

lemma *i-set-Minus-closed*:
 [$I \in i\text{-set}; i\text{Min } I \leq k$] $\implies k \ominus I \in i\text{-set}$
 ⟨proof⟩

lemma *i-set0-Plus-neg-closed*: $I \in i\text{-set0} \implies I \oplus -k \in i\text{-set0}$
 ⟨proof⟩

lemma *i-set0-Minus-closed*: $I \in i\text{-set0} \implies k \ominus I \in i\text{-set0}$
 ⟨proof⟩

lemmas *i-set-IntOp-closed* =
i-set-Plus-closed
i-set-Mult-closed
i-set-Plus-neg-closed
i-set-Minus-closed

lemmas *i-set0-IntOp-closed* =
i-set0-Plus-closed
i-set0-Mult-closed
i-set0-Plus-neg-closed
i-set0-Minus-closed

2.5.5 Interval sets are not closed under division

lemma *iMOD-div-mod-gr0-not-in-i-set*:
 $\llbracket 0 < k; k < m; 0 < m \bmod k \rrbracket \implies [r, \bmod m] \odot k \notin i\text{-set}$
<proof>

lemma *iMODb-div-mod-gr0-not-in-i-set*:
 $\llbracket 0 < k; k < m; 0 < m \bmod k; k \leq c \rrbracket \implies [r, \bmod m, c] \odot k \notin i\text{-set}$
<proof>

lemma $[0, \bmod 3] \odot 2 \notin i\text{-set}$
<proof>

lemma *i-set-Div-not-closed*: $\text{Suc } 0 < k \implies \exists I \in i\text{-set}. I \odot k \notin i\text{-set}$
<proof>

lemma *i-set0-Div-not-closed*: $\text{Suc } 0 < k \implies \exists I \in i\text{-set0}. I \odot k \notin i\text{-set0}$
<proof>

2.5.6 Sets of intervals closed under division

inductive-set *NatMultiples* :: *nat set* \Rightarrow *nat set*
for *F* :: *nat set*

where

NatMultiples-Factor: $k \in F \implies k \in \text{NatMultiples } F$

NatMultiples-Product: $\llbracket k \in F; m \in \text{NatMultiples } F \rrbracket \implies k * m \in \text{NatMultiples } F$

lemma *NatMultiples-ex-divisor*: $m \in \text{NatMultiples } F \implies \exists k \in F. m \bmod k = 0$
<proof>

lemma *NatMultiples-product-factor*: $\llbracket a \in F; b \in F \rrbracket \implies a * b \in \text{NatMultiples } F$
<proof>

lemma *NatMultiples-product-factor-multiple*:
 $\llbracket a \in F; b \in \text{NatMultiples } F \rrbracket \implies a * b \in \text{NatMultiples } F$
<proof>

lemma *NatMultiples-product-multiple-factor*:
 $\llbracket a \in \text{NatMultiples } F; b \in F \rrbracket \implies a * b \in \text{NatMultiples } F$
<proof>

lemma *NatMultiples-product-multiple*:
 $\llbracket a \in \text{NatMultiples } F; b \in \text{NatMultiples } F \rrbracket \implies a * b \in \text{NatMultiples } F$
<proof>

Subset of $i\text{-set}$ containing only continuous intervals, i. e., without $iMOD$ and $iMODb$.

inductive-set $i\text{-set-cont} :: (\text{nat set}) \text{ set}$
where

$i\text{-set-cont-FROM}[\text{intro}]: [n..] \in i\text{-set-cont}$
 $| i\text{-set-cont-TILL}[\text{intro}]: [...n] \in i\text{-set-cont}$
 $| i\text{-set-cont-IN}[\text{intro}]: [n..,d] \in i\text{-set-cont}$

definition $i\text{-set0-cont} :: (\text{nat set}) \text{ set}$

where $i\text{-set0-cont} \equiv \text{insert } \{\} i\text{-set-cont}$

lemma $i\text{-set-cont-subset-i-set}: i\text{-set-cont} \subseteq i\text{-set}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-cont-subset-i-set0}: i\text{-set0-cont} \subseteq i\text{-set0}$
 $\langle \text{proof} \rangle$

Minimal definition of $i\text{-set-cont}$

inductive-set $i\text{-set-cont-min} :: (\text{nat set}) \text{ set}$
where

$i\text{-set-cont-FROM}[\text{intro}]: [n..] \in i\text{-set-cont-min}$
 $| i\text{-set-cont-IN}[\text{intro}]: [n..,d] \in i\text{-set-cont-min}$

definition $i\text{-set0-cont-min} :: (\text{nat set}) \text{ set}$

where $i\text{-set0-cont-min} \equiv \text{insert } \{\} i\text{-set-cont-min}$

lemma $i\text{-set-cont-min-eq}: i\text{-set-cont} = i\text{-set-cont-min}$
 $\langle \text{proof} \rangle$

$i\text{next}$ and $i\text{prev}$ with continuous intervals

lemma $i\text{-set-cont-i\text{next}}$:

$\llbracket I \in i\text{-set-cont}; n \in I; \text{finite } I \implies n < \text{Max } I \rrbracket \implies i\text{next } n \ I = \text{Suc } n$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-cont-i\text{prev}}$:

$\llbracket I \in i\text{-set-cont}; n \in I; i\text{Min } I < n \rrbracket \implies i\text{prev } n \ I = n - \text{Suc } 0$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-cont-i\text{next--less}}$:

$\llbracket I \in i\text{-set-cont}; n \in I; n < n0; n0 \in I \rrbracket \implies i\text{next } n \ I = \text{Suc } n$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-cont-i\text{prev--greater}}$:

$\llbracket I \in i\text{-set-cont}; n \in I; n0 < n; n0 \in I \rrbracket \implies i\text{prev } n \ I = n - \text{Suc } 0$
 $\langle \text{proof} \rangle$

Sets of modulo intervals

inductive-set $i\text{-set-mult} :: \text{nat} \Rightarrow (\text{nat set}) \text{ set}$

for $k :: \text{nat}$
where
 $i\text{-set-mult-FROM}[intro!]: [n..] \in i\text{-set-mult } k$
 $| i\text{-set-mult-TILL}[intro!]: [..n] \in i\text{-set-mult } k$
 $| i\text{-set-mult-IN}[intro!]: [n..,d] \in i\text{-set-mult } k$
 $| i\text{-set-mult-MOD}[intro!]: [r, \text{mod } m * k] \in i\text{-set-mult } k$
 $| i\text{-set-mult-MODb}[intro!]: [r, \text{mod } m * k, c] \in i\text{-set-mult } k$

definition $i\text{-set0-mult} :: \text{nat} \Rightarrow (\text{nat set}) \text{ set}$
where $i\text{-set0-mult } k \equiv \text{insert } \{\} (i\text{-set-mult } k)$

lemma
 $i\text{-set0-mult-empty}[intro!]: \{\} \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-FROM}[intro!]: [n..] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-TILL}[intro!]: [..n] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-IN}[intro!]: [n..,d] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-MOD}[intro!]: [r, \text{mod } m * k] \in i\text{-set0-mult } k$ **and**
 $i\text{-set0-mult-MODb}[intro!]: [r, \text{mod } m * k, c] \in i\text{-set0-mult } k$
 $\langle \text{proof} \rangle$

lemmas $i\text{-set0-mult-intros} =$
 $i\text{-set0-mult-empty}$
 $i\text{-set0-mult-FROM}$
 $i\text{-set0-mult-TILL}$
 $i\text{-set0-mult-IN}$
 $i\text{-set0-mult-MOD}$
 $i\text{-set0-mult-MODb}$

lemma $i\text{-set-mult-subset-i-set0-mult}: i\text{-set-mult } k \subseteq i\text{-set0-mult } k$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-mult-subset-i-set}: i\text{-set-mult } k \subseteq i\text{-set}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-mult-subset-i-set0}: i\text{-set0-mult } k \subseteq i\text{-set0}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-mult-0-eq-i-set-cont}: i\text{-set-mult } 0 = i\text{-set-cont}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-mult-0-eq-i-set0-cont}: i\text{-set0-mult } 0 = i\text{-set0-cont}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set-mult-1-eq-i-set}: i\text{-set-mult } (\text{Suc } 0) = i\text{-set}$
 $\langle \text{proof} \rangle$

lemma $i\text{-set0-mult-1-eq-i-set0}: i\text{-set0-mult } (\text{Suc } 0) = i\text{-set0}$
 $\langle \text{proof} \rangle$

lemma *i-set-mult-imp-not-empty*: $I \in i\text{-set-mult } k \implies I \neq \{\}$
 ⟨proof⟩

lemma *iMOD-in-i-set-mult-imp-divisor-mod-0*:
 $\llbracket m \neq \text{Suc } 0; [r, \text{mod } m] \in i\text{-set-mult } k \rrbracket \implies m \text{ mod } k = 0$
 ⟨proof⟩

lemma
divisor-mod-0-imp-iMOD-in-i-set-mult: $m \text{ mod } k = 0 \implies [r, \text{mod } m] \in i\text{-set-mult } k$ **and**
divisor-mod-0-imp-iMODb-in-i-set-mult: $m \text{ mod } k = 0 \implies [r, \text{mod } m, c] \in i\text{-set-mult } k$
 ⟨proof⟩

lemma *iMOD-in-i-set-mult--divisor-mod-0-conv*:
 $m \neq \text{Suc } 0 \implies ([r, \text{mod } m] \in i\text{-set-mult } k) = (m \text{ mod } k = 0)$
 ⟨proof⟩

lemma *i-set-mult-neq1-subset-i-set*: $k \neq \text{Suc } 0 \implies i\text{-set-mult } k \subset i\text{-set}$
 ⟨proof⟩

lemma *mod-0-imp-i-set-mult-subset*:
 $a \text{ mod } b = 0 \implies i\text{-set-mult } a \subseteq i\text{-set-mult } b$
 ⟨proof⟩

lemma *i-set-mult-subset-imp-mod-0*:
 $\llbracket a \neq \text{Suc } 0; (i\text{-set-mult } a \subseteq i\text{-set-mult } b) \rrbracket \implies a \text{ mod } b = 0$
 ⟨proof⟩

lemma *i-set-mult-subset-conv*:
 $a \neq \text{Suc } 0 \implies (i\text{-set-mult } a \subseteq i\text{-set-mult } b) = (a \text{ mod } b = 0)$
 ⟨proof⟩

lemma *i-set-mult-mod-0-div*:
 $\llbracket I \in i\text{-set-mult } k; k \text{ mod } d = 0 \rrbracket \implies I \odot d \in i\text{-set-mult } (k \text{ div } d)$
 ⟨proof⟩

Intervals from *i-set-mult* k remain in *i-set* after division by d a divisor of k .

corollary *i-set-mult-mod-0-div-i-set*:
 $\llbracket I \in i\text{-set-mult } k; k \text{ mod } d = 0 \rrbracket \implies I \odot d \in i\text{-set}$
 ⟨proof⟩

corollary *i-set-mult-div-self-i-set*:
 $I \in i\text{-set-mult } k \implies I \odot k \in i\text{-set}$
 ⟨proof⟩

lemma *i-set-mod-0-mult-in-i-set-mult*:
 $\llbracket I \in i\text{-set}; m \text{ mod } k = 0 \rrbracket \implies I \otimes m \in i\text{-set-mult } k$

<proof>

lemma *i-set-self-mult-in-i-set-mult*:

$I \in i\text{-set} \implies I \otimes k \in i\text{-set-mult } k$

<proof>

lemma *i-set-mult-mod-gr0-div-not-in-i-set*:

$\llbracket 0 < k; 0 < d; 0 < k \bmod d \rrbracket \implies \exists I \in i\text{-set-mult } k. I \odot d \notin i\text{-set}$

<proof>

lemma *i-set0-mult-mod-0-div*:

$\llbracket I \in i\text{-set0-mult } k; k \bmod d = 0 \rrbracket \implies I \odot d \in i\text{-set0-mult } (k \text{ div } d)$

<proof>

corollary *i-set0-mult-mod-0-div-i-set0*:

$\llbracket I \in i\text{-set0-mult } k; k \bmod d = 0 \rrbracket \implies I \odot d \in i\text{-set0}$

<proof>

corollary *i-set0-mult-div-self-i-set0*:

$I \in i\text{-set0-mult } k \implies I \odot k \in i\text{-set0}$

<proof>

lemma *i-set0-mod-0-mult-in-i-set0-mult*:

$\llbracket I \in i\text{-set0}; m \bmod k = 0 \rrbracket \implies I \otimes m \in i\text{-set0-mult } k$

<proof>

lemma *i-set0-self-mult-in-i-set0-mult*:

$I \in i\text{-set0} \implies I \otimes k \in i\text{-set0-mult } k$

<proof>

lemma *i-set0-mult-mod-gr0-div-not-in-i-set0*:

$\llbracket 0 < k; 0 < d; 0 < k \bmod d \rrbracket \implies \exists I \in i\text{-set0-mult } k. I \odot d \notin i\text{-set0}$

<proof>

end

3 Temporal logic operators on natural intervals

theory *IL-TemporalOperators*

imports *IL-IntervalOperators*

begin

Bool : some additional properties

instantiation *bool* :: {ord, zero, one, plus, times, order}

begin

definition *Zero-bool-def* [simp]: $0 \equiv \text{False}$

definition *One-bool-def* [simp]: $1 \equiv \text{True}$

definition *add-bool-def*: $a + b \equiv a \vee b$

definition *mult-bool-def*: $a * b \equiv a \wedge b$

instance $\langle proof \rangle$

end

value $False < True$

value $True < True$

value $True \leq True$

lemmas *bool-op-rel-defs* =

add-bool-def

mult-bool-def

less-bool-def

le-bool-def

instance *bool* :: *wellorder*

$\langle proof \rangle$

instance *bool* :: *comm-semiring-1*

$\langle proof \rangle$

3.1 Basic definitions

lemma *UNIV-nat*: $\mathbf{N} = (UNIV::nat\ set)$

$\langle proof \rangle$

Universal temporal operator: Always/Globally

definition *iAll* :: $iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Always

where $iAll\ I\ P \equiv \forall t \in I. P\ t$

Existential temporal operator: Eventually/Finally

definition *iEx* :: $iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Eventually

where $iEx\ I\ P \equiv \exists t \in I. P\ t$

syntax

$-iAll :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \square - \cdot / -)\ [0, 0, 10]\ 10)$

$-iEx :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \diamond - \cdot / -)\ [0, 0, 10]\ 10)$

translations

$\square\ t\ I. P \equiv CONST\ iAll\ I\ (\lambda t. P)$

$\diamond\ t\ I. P \equiv CONST\ iEx\ I\ (\lambda t. P)$

Future temporal operator: Next

definition *iNext* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Next

where $iNext\ t0\ I\ P \equiv P\ (inext\ t0\ I)$

Past temporal operator: Last/Previous

definition *iLast* :: $Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Last

where $iLast\ t0\ I\ P \equiv P\ (iprev\ t0\ I)$

syntax

$-iNext :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \circ - - - / -)\ [0, 0, 10]\ 10)$

$-iLast :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \oplus - - - / -)\ [0, 0, 10]\ 10)$

translations

$\circ\ t\ t0\ I.\ P \equiv CONST\ iNext\ t0\ I\ (\lambda t.\ P)$

$\ominus\ t\ t0\ I.\ P \equiv CONST\ iLast\ t0\ I\ (\lambda t.\ P)$

lemma $\circ\ t\ 10\ [0\dots].\ (t + 10 > 10)$

$\langle proof \rangle$

The following versions of Next and Last operator differ in the cases where no next/previous element exists or specified time point is not in interval: the weak versions return *True* and the strong versions return *False*.

definition $iNextWeak :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Weak Next
where $iNextWeak\ t0\ I\ P \equiv (\square\ t\ \{inext\ t0\ I\}\ \downarrow >\ t0.\ P\ t)$

definition $iNextStrong :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Strong Next
where $iNextStrong\ t0\ I\ P \equiv (\diamond\ t\ \{inext\ t0\ I\}\ \downarrow >\ t0.\ P\ t)$

definition $iLastWeak :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Weak Last
where $iLastWeak\ t0\ I\ P \equiv (\square\ t\ \{iprev\ t0\ I\}\ \downarrow <\ t0.\ P\ t)$

definition $iLastStrong :: Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Strong Last
where $iLastStrong\ t0\ I\ P \equiv (\diamond\ t\ \{iprev\ t0\ I\}\ \downarrow <\ t0.\ P\ t)$

syntax

$-iNextWeak :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \circ_W - - - / -)\ [0, 0, 10]\ 10)$

$-iNextStrong :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \circ_S - - - / -)\ [0, 0, 10]\ 10)$

$-iLastWeak :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \oplus_W - - - / -)\ [0, 0, 10]\ 10)$

$-iLastStrong :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow bool\ ((\exists \oplus_S - - - / -)\ [0, 0, 10]\ 10)$

translations

$\circ_W\ t\ t0\ I.\ P \equiv CONST\ iNextWeak\ t0\ I\ (\lambda t.\ P)$

$\circ_S\ t\ t0\ I.\ P \equiv CONST\ iNextStrong\ t0\ I\ (\lambda t.\ P)$

$\ominus_W\ t\ t0\ I.\ P \equiv CONST\ iLastWeak\ t0\ I\ (\lambda t.\ P)$

$\ominus_S\ t\ t0\ I.\ P \equiv CONST\ iLastStrong\ t0\ I\ (\lambda t.\ P)$

Some examples for Next and Last operator

lemma $\circ\ t\ 5\ [0\dots,10].\ ([0::int,10,20,30,40,50,60,70,80,90] ! t < 80)$

$\langle proof \rangle$

lemma $\ominus t 5 [0\dots,10]. ([0::int,10,20,30,40,50,60,70,80,90] ! t < 80)$
 ⟨proof⟩

Temporal Until operator

definition $iUntil :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Until
where $iUntil I P Q \equiv \diamond t I. Q t \wedge (\Box t' (I \downarrow < t). P t')$

Temporal Since operator (past operator corresponding to Until)

definition $iSince :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ — Since
where $iSince I P Q \equiv \diamond t I. Q t \wedge (\Box t' (I \downarrow > t). P t')$

syntax

$-iUntil :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (3U - -)./ -) [10, 0, 0, 0, 10] 10)$
 $-iSince :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (3S - -)./ -) [10, 0, 0, 0, 10] 10)$

translations

$P. t U t' I. Q \equiv CONST iUntil I (\lambda t. P) (\lambda t'. Q)$
 $P. t S t' I. Q \equiv CONST iSince I (\lambda t. P) (\lambda t'. Q)$

definition $iWeakUntil :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
 Weak Until/Wating-for/Unless

where $iWeakUntil I P Q \equiv (\Box t I. P t) \vee (\diamond t I. Q t \wedge (\Box t' (I \downarrow < t). P t'))$

definition $iWeakSince :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
 Weak Since/Back-to

where $iWeakSince I P Q \equiv (\Box t I. P t) \vee (\diamond t I. Q t \wedge (\Box t' (I \downarrow > t). P t'))$

syntax

$-iWeakUntil :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (3W - -)./ -) [10, 0, 0, 0, 10] 10)$
 $-iWeakSince :: Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (3B - -)./ -) [10, 0, 0, 0, 10] 10)$

translations

$P. t W t' I. Q \equiv CONST iWeakUntil I (\lambda t. P) (\lambda t'. Q)$
 $P. t B t' I. Q \equiv CONST iWeakSince I (\lambda t. P) (\lambda t'. Q)$

definition $iRelease :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
 Release

where $iRelease I P Q \equiv (\Box t I. Q t) \vee (\diamond t I. P t \wedge (\Box t' (I \downarrow \leq t). Q t'))$

definition $iTrigger :: iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$ —
 Trigger

where $iTrigger I P Q \equiv (\Box t I. Q t) \vee (\diamond t I. P t \wedge (\Box t' (I \downarrow \geq t). Q t'))$

syntax

-iRelease :: $Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (\mathcal{R} - -)./) -) [10, 0, 0, 0, 10] 10$
-iTrigger :: $Time \Rightarrow Time \Rightarrow iT \Rightarrow (Time \Rightarrow bool) \Rightarrow (Time \Rightarrow bool) \Rightarrow bool$
 $((./ - (\mathcal{T} - -)./) -) [10, 0, 0, 0, 10] 10$

translations

$P. t \mathcal{R} t' I. Q \equiv \text{CONST } i\text{Release } I (\lambda t. P) (\lambda t'. Q)$
 $P. t \mathcal{T} t' I. Q \equiv \text{CONST } i\text{Trigger } I (\lambda t. P) (\lambda t'. Q)$

lemmas *iTL-Next-defs* =

iNext-def *iLast-def*
iNextWeak-def *iLastWeak-def*
iNextStrong-def *iLastStrong-def*

lemmas *iTL-defs* =

iAll-def *iEx-def*
iUntil-def *iSince-def*
iWeakUntil-def *iWeakSince-def*
iRelease-def *iTrigger-def*
iTL-Next-defs

$\langle ML \rangle$

3.2 Basic lemmata for temporal operators

3.2.1 Intro/elim rules

lemma

iexI[intro]: $\llbracket P t; t \in I \rrbracket \Longrightarrow \diamond t I. P t$ **and**
rev-iexI[intro?]: $\llbracket t \in I; P t \rrbracket \Longrightarrow \diamond t I. P t$ **and**
iexE[elim!]: $\llbracket \diamond t I. P t; \bigwedge t. \llbracket t \in I; P t \rrbracket \Longrightarrow Q \rrbracket \Longrightarrow Q$
 $\langle \text{proof} \rangle$

lemma

iallI[intro!]: $(\bigwedge t. t \in I \Longrightarrow P t) \Longrightarrow \square t I. P t$ **and**
ispec[dest?]: $\llbracket \square t I. P t; t \in I \rrbracket \Longrightarrow P t$ **and**
iallE[elim]: $\llbracket \square t I. P t; P t \Longrightarrow Q; t \notin I \Longrightarrow Q \rrbracket \Longrightarrow Q$
 $\langle \text{proof} \rangle$

lemma

iuntilI[intro]:
 $\llbracket Q t; (\bigwedge t'. t' \in I \downarrow < t \Longrightarrow P t'); t \in I \rrbracket \Longrightarrow P t'. t' \mathcal{U} t I. Q t$ **and**
rev-iuntilI[intro?]:
 $\llbracket t \in I; Q t; (\bigwedge t'. t' \in I \downarrow < t \Longrightarrow P t') \rrbracket \Longrightarrow P t'. t' \mathcal{U} t I. Q t$
 $\langle \text{proof} \rangle$

lemma

iuntilE[elim]:
 $\llbracket P' t'. t' \mathcal{U} t I. P t; \bigwedge t. \llbracket t \in I; P t \rrbracket \Longrightarrow Q \rrbracket \Longrightarrow Q$
 $\langle \text{proof} \rangle$

lemma

isinceI[intro]:

$\llbracket Q\ t; (\bigwedge t'. t' \in I \downarrow > t \implies P\ t'); t \in I \rrbracket \implies P\ t'. t' \mathcal{S}\ t\ I.\ Q\ t$ **and**

rev-isinceI[intro?]:

$\llbracket t \in I; Q\ t; (\bigwedge t'. t' \in I \downarrow > t \implies P\ t') \rrbracket \implies P\ t'. t' \mathcal{S}\ t\ I.\ Q\ t$

<proof>

lemma

isinceE[elim]:

$\llbracket P'\ t'. t' \mathcal{S}\ t\ I.\ P\ t; \bigwedge t. \llbracket t \in I; P\ t \rrbracket \implies Q \rrbracket \implies Q$

<proof>

3.2.2 Rewrite rules for trivial simplification

lemma *iall-triv[simp]:* $(\Box\ t\ I.\ P) = ((\exists t. t \in I) \longrightarrow P)$

<proof>

lemma *isex-triv[simp]:* $(\Diamond\ t\ I.\ P) = ((\exists t. t \in I) \wedge P)$

<proof>

lemma *isex-conjL1:*

$(\Diamond\ t1\ I1.\ (P1\ t1 \wedge (\Diamond\ t2\ I2.\ P2\ t1\ t2))) =$

$(\Diamond\ t1\ I1.\ \Diamond\ t2\ I2.\ P1\ t1 \wedge P2\ t1\ t2)$

<proof>

lemma *isex-conjR1:*

$(\Diamond\ t1\ I1.\ ((\Diamond\ t2\ I2.\ P2\ t1\ t2) \wedge P1\ t1)) =$

$(\Diamond\ t1\ I1.\ \Diamond\ t2\ I2.\ P2\ t1\ t2 \wedge P1\ t1)$

<proof>

lemma *isex-conjL2:*

$(\Diamond\ t1\ I1.\ (P1\ t1 \wedge (\Diamond\ t2\ (I2\ t1).\ P2\ t1\ t2))) =$

$(\Diamond\ t1\ I1.\ \Diamond\ t2\ (I2\ t1).\ P1\ t1 \wedge P2\ t1\ t2)$

<proof>

lemma *isex-conjR2:*

$(\Diamond\ t1\ I1.\ ((\Diamond\ t2\ (I2\ t1).\ P2\ t1\ t2) \wedge P1\ t1)) =$

$(\Diamond\ t1\ I1.\ \Diamond\ t2\ (I2\ t1).\ P2\ t1\ t2 \wedge P1\ t1)$

<proof>

lemma *isex-commute:*

$(\Diamond\ t1\ I1.\ \Diamond\ t2\ I2.\ P\ t1\ t2) =$

$(\Diamond\ t2\ I2.\ \Diamond\ t1\ I1.\ P\ t1\ t2)$

<proof>

lemma *iall-conjL1:*

$I2 \neq \{\} \implies$

$(\Box\ t1\ I1.\ (P1\ t1 \wedge (\Box\ t2\ I2.\ P2\ t1\ t2))) =$

$(\Box\ t1\ I1.\ \Box\ t2\ I2.\ P1\ t1 \wedge P2\ t1\ t2)$

$\langle proof \rangle$

lemma *iall-conjR1*:

$$\begin{aligned} I2 \neq \{\} &\implies \\ (\Box t1 I1. ((\Box t2 I2. P2 t1 t2) \wedge P1 t1)) &= \\ (\Box t1 I1. \Box t2 I2. P2 t1 t2 \wedge P1 t1) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iall-conjL2*:

$$\begin{aligned} \Box t1 I1. I2 t1 \neq \{\} &\implies \\ (\Box t1 I1. (P1 t1 \wedge (\Box t2 (I2 t1). P2 t1 t2))) &= \\ (\Box t1 I1. \Box t2 (I2 t1). P1 t1 \wedge P2 t1 t2) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iall-conjR2*:

$$\begin{aligned} \Box t1 I1. I2 t1 \neq \{\} &\implies \\ (\Box t1 I1. ((\Box t2 (I2 t1). P2 t1 t2) \wedge P1 t1)) &= \\ (\Box t1 I1. \Box t2 (I2 t1). P2 t1 t2 \wedge P1 t1) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iall-commute*:

$$\begin{aligned} (\Box t1 I1. \Box t2 I2. P t1 t2) &= \\ (\Box t2 I2. \Box t1 I1. P t1 t2) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iall-conj-distrib*:

$$\begin{aligned} (\Box t I. P t \wedge Q t) &= ((\Box t I. P t) \wedge (\Box t I. Q t)) \\ \langle proof \rangle & \end{aligned}$$

lemma *iox-disj-distrib*:

$$\begin{aligned} (\Diamond t I. P t \vee Q t) &= ((\Diamond t I. P t) \vee (\Diamond t I. Q t)) \\ \langle proof \rangle & \end{aligned}$$

lemma *iall-conj-distrib2*:

$$\begin{aligned} (\Box t1 I1. \Box t2 (I2 t1). P t1 t2 \wedge Q t1 t2) &= \\ ((\Box t1 I1. \Box t2 (I2 t1). P t1 t2) \wedge (\Box t1 I1. \Box t2 (I2 t1). Q t1 t2)) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iox-disj-distrib2*:

$$\begin{aligned} (\Diamond t1 I1. \Diamond t2 (I2 t1). P t1 t2 \vee Q t1 t2) &= \\ ((\Diamond t1 I1. \Diamond t2 (I2 t1). P t1 t2) \vee (\Diamond t1 I1. \Diamond t2 (I2 t1). Q t1 t2)) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iUntil-disj-distrib*:

$$\begin{aligned} (P t1. t1 \mathcal{U} t2 I. (Q1 t2 \vee Q2 t2)) &= ((P t1. t1 \mathcal{U} t2 I. Q1 t2) \vee (P t1. t1 \mathcal{U} \\ t2 I. Q2 t2)) & \\ \langle proof \rangle & \end{aligned}$$

lemma *iSince-disj-distrib*:

$$(P \ t1. \ t1 \ S \ t2 \ I. \ (Q1 \ t2 \ \vee \ Q2 \ t2)) = ((P \ t1. \ t1 \ S \ t2 \ I. \ Q1 \ t2) \ \vee \ (P \ t1. \ t1 \ S \ t2 \ I. \ Q2 \ t2))$$

<proof>

lemma

$$iNext\text{-iff}: (\circ \ t \ t0 \ I. \ P \ t) = (\square \ t \ [\dots 0] \oplus (inext \ t0 \ I). \ P \ t) \ \mathbf{and}$$

$$iLast\text{-iff}: (\ominus \ t \ t0 \ I. \ P \ t) = (\square \ t \ [\dots 0] \oplus (iprev \ t0 \ I). \ P \ t)$$

<proof>

lemma

$$iNext\text{-iEx-iff}: (\circ \ t \ t0 \ I. \ P \ t) = (\diamond \ t \ [\dots 0] \oplus (inext \ t0 \ I). \ P \ t) \ \mathbf{and}$$

$$iLast\text{-iEx-iff}: (\ominus \ t \ t0 \ I. \ P \ t) = (\diamond \ t \ [\dots 0] \oplus (iprev \ t0 \ I). \ P \ t)$$

<proof>

lemma *inext-singleton-cut-greater-not-empty-iff:*

$$(\{inext \ t0 \ I\} \ \downarrow > \ t0 \ \neq \ \{\}) = (I \ \downarrow > \ t0 \ \neq \ \{\} \ \wedge \ t0 \ \in \ I)$$

<proof>

lemma *iprev-singleton-cut-less-not-empty-iff:*

$$(\{iprev \ t0 \ I\} \ \downarrow < \ t0 \ \neq \ \{\}) = (I \ \downarrow < \ t0 \ \neq \ \{\} \ \wedge \ t0 \ \in \ I)$$

<proof>

lemma *inext-singleton-cut-greater-empty-iff:*

$$(\{inext \ t0 \ I\} \ \downarrow > \ t0 \ = \ \{\}) = (I \ \downarrow > \ t0 \ = \ \{\} \ \vee \ t0 \ \notin \ I)$$

<proof>

lemma *iprev-singleton-cut-less-empty-iff:*

$$(\{iprev \ t0 \ I\} \ \downarrow < \ t0 \ = \ \{\}) = (I \ \downarrow < \ t0 \ = \ \{\} \ \vee \ t0 \ \notin \ I)$$

<proof>

lemma *iNextWeak-iff* : $(\circ_W \ t \ t0 \ I. \ P \ t) = ((\circ \ t \ t0 \ I. \ P \ t) \ \vee \ (I \ \downarrow > \ t0 \ = \ \{\}) \ \vee \ t0 \ \notin \ I)$

<proof>

lemma *iNextStrong-iff* : $(\circ_S \ t \ t0 \ I. \ P \ t) = ((\circ \ t \ t0 \ I. \ P \ t) \ \wedge \ (I \ \downarrow > \ t0 \ \neq \ \{\}) \ \wedge \ t0 \ \in \ I)$

<proof>

lemma *iLastWeak-iff* : $(\ominus_W \ t \ t0 \ I. \ P \ t) = ((\ominus \ t \ t0 \ I. \ P \ t) \ \vee \ (I \ \downarrow < \ t0 \ = \ \{\}) \ \vee \ t0 \ \notin \ I)$

<proof>

lemma *iLastStrong-iff* : $(\ominus_S \ t \ t0 \ I. \ P \ t) = ((\ominus \ t \ t0 \ I. \ P \ t) \ \wedge \ (I \ \downarrow < \ t0 \ \neq \ \{\}) \ \wedge \ t0 \ \in \ I)$

<proof>

lemmas *iTL-Next-iff* =

$$iNext\text{-iff} \ iLast\text{-iff}$$

iNextWeak-iff iNextStrong-iff
iLastWeak-iff iLastStrong-iff

lemma

iNext-iff-singleton : $(\circlearrowleft t \ t0 \ I. \ P \ t) = (\square t \ \{inext \ t0 \ I\}. \ P \ t)$ **and**
iLast-iff-singleton : $(\ominus t \ t0 \ I. \ P \ t) = (\square t \ \{iprev \ t0 \ I\}. \ P \ t)$

<proof>

lemmas *iNextLast-iff-singleton* = *iNext-iff-singleton iLast-iff-singleton*

lemma

iNext-iEx-iff-singleton : $(\circlearrowleft t \ t0 \ I. \ P \ t) = (\diamond t \ \{inext \ t0 \ I\}. \ P \ t)$ **and**
iLast-iEx-iff-singleton : $(\ominus t \ t0 \ I. \ P \ t) = (\diamond t \ \{iprev \ t0 \ I\}. \ P \ t)$

<proof>

lemma

iNextWeak-iAll-conv: $(\circlearrowleft_W t \ t0 \ I. \ P \ t) = (\square t \ (\{inext \ t0 \ I\} \ \downarrow > \ t0). \ P \ t)$ **and**
iNextStrong-iEx-conv: $(\circlearrowleft_S t \ t0 \ I. \ P \ t) = (\diamond t \ (\{inext \ t0 \ I\} \ \downarrow > \ t0). \ P \ t)$ **and**
iLastWeak-iAll-conv: $(\ominus_W t \ t0 \ I. \ P \ t) = (\square t \ (\{iprev \ t0 \ I\} \ \downarrow < \ t0). \ P \ t)$ **and**
iLastStrong-iEx-conv: $(\ominus_S t \ t0 \ I. \ P \ t) = (\diamond t \ (\{iprev \ t0 \ I\} \ \downarrow < \ t0). \ P \ t)$

<proof>

lemma

iAll-True[simp]: $\square t \ I. \ True$ **and**
iAll-False[simp]: $(\square t \ I. \ False) = (I = \{\})$ **and**
iEx-True[simp]: $(\diamond t \ I. \ True) = (I \neq \{\})$ **and**
iEx-False[simp]: $\neg (\diamond t \ I. \ False)$

<proof>

lemma *empty-iff-iAll-False*: $(I = \{\}) = (\square t \ I. \ False)$ *<proof>*

lemma *not-empty-iff-iEx-True*: $(I \neq \{\}) = (\diamond t \ I. \ True)$ *<proof>*

lemma

iNext-True: $\circlearrowleft t \ t0 \ I. \ True$ **and**
iNextWeak-True: $(\circlearrowleft_W t \ t0 \ I. \ True)$ **and**
iNext-False: $\neg (\circlearrowleft t \ t0 \ I. \ False)$ **and**
iNextStrong-False: $\neg (\circlearrowleft_S t \ t0 \ I. \ False)$

<proof>

lemma

iNextStrong-True: $(\circlearrowleft_S t \ t0 \ I. \ True) = (I \ \downarrow > \ t0 \ \neq \ \{\} \ \wedge \ t0 \ \in \ I)$ **and**
iNextWeak-False: $(\neg (\circlearrowleft_W t \ t0 \ I. \ False)) = (I \ \downarrow > \ t0 \ \neq \ \{\} \ \wedge \ t0 \ \in \ I)$

<proof>

lemma

iLast-True: $\ominus t \ t0 \ I. \ True$ **and**

iLastWeak-True: $(\ominus_W t t0 I. True)$ **and**
iLast-False: $\neg (\ominus t t0 I. False)$ **and**
iLastStrong-False: $\neg (\ominus_S t t0 I. False)$
 ⟨proof⟩

lemma

iLastStrong-True: $(\ominus_S t t0 I. True) = (I \downarrow < t0 \neq \{\} \wedge t0 \in I)$ **and**
iLastWeak-False: $(\neg (\ominus_W t t0 I. False)) = (I \downarrow < t0 \neq \{\} \wedge t0 \in I)$
 ⟨proof⟩

lemma *iUntil-True-left[simp]*: $(True. t' U t I. Q t) = (\diamond t I. Q t)$
 ⟨proof⟩

lemma *iUntil-True[simp]*: $(P t'. t' U t I. True) = (I \neq \{\})$
 ⟨proof⟩

lemma *iUntil-False-left[simp]*: $(False. t' U t I. Q t) = (I \neq \{\} \wedge Q (iMin I))$
 ⟨proof⟩

lemma *iUntil-False[simp]*: $\neg (P t'. t' U t I. False)$
 ⟨proof⟩

lemma *iSince-True-left[simp]*: $(True. t' S t I. Q t) = (\diamond t I. Q t)$
 ⟨proof⟩

lemma *iSince-True-if*:
 $(P t'. t' S t I. True) =$
(if finite I then I ≠ {} else $\diamond t1 I. \square t2 (I \downarrow > t1). P t2$)
 ⟨proof⟩

corollary *iSince-True-finite[simp]*: $finite I \implies (P t'. t' S t I. True) = (I \neq \{\})$
 ⟨proof⟩

lemma *iSince-False-left[simp]*: $(False. t' S t I. Q t) = (finite I \wedge I \neq \{\} \wedge Q (Max I))$
 ⟨proof⟩

lemma *iSince-False[simp]*: $\neg (P t'. t' S t I. False)$
 ⟨proof⟩

lemma *iWeakUntil-True-left[simp]*: $True. t' W t I. Q t$
 ⟨proof⟩

lemma *iWeakUntil-True[simp]*: $P t'. t' W t I. True$
 ⟨proof⟩

lemma *iWeakUntil-False-left[simp]*: $(False. t' W t I. Q t) = (I = \{\} \vee Q (iMin I))$
 ⟨proof⟩

lemma *iWeakUntil-False[simp]*: $(P\ t'.\ t'\ \mathcal{W}\ t\ I.\ \text{False}) = (\Box\ t\ I.\ P\ t)$
 $\langle\text{proof}\rangle$

lemma *iWeakSince-True-left[simp]*: $\text{True}.\ t'\ \mathcal{B}\ t\ I.\ Q\ t$
 $\langle\text{proof}\rangle$

lemma *iWeakSince-True-disj*:
 $(P\ t'.\ t'\ \mathcal{B}\ t\ I.\ \text{True}) =$
 $(I = \{\}) \vee (\Diamond\ t1\ I.\ \Box\ t2\ (I\ \Downarrow\ t1).\ P\ t2))$
 $\langle\text{proof}\rangle$

lemma *iWeakSince-True-finite[simp]*: $\text{finite}\ I \implies (P\ t'.\ t'\ \mathcal{B}\ t\ I.\ \text{True})$
 $\langle\text{proof}\rangle$

lemma *iWeakSince-False-left[simp]*: $(\text{False}.\ t'\ \mathcal{B}\ t\ I.\ Q\ t) = (I = \{\}) \vee (\text{finite}\ I \wedge$
 $Q\ (\text{Max}\ I))$
 $\langle\text{proof}\rangle$

lemma *iWeakSince-False[simp]*: $(P\ t'.\ t'\ \mathcal{B}\ t\ I.\ \text{False}) = (\Box\ t\ I.\ P\ t)$
 $\langle\text{proof}\rangle$

lemma *iRelease-True-left[simp]*: $(\text{True}.\ t'\ \mathcal{R}\ t\ I.\ Q\ t) = (I = \{\}) \vee Q\ (\text{iMin}\ I)$
 $\langle\text{proof}\rangle$

lemma *iRelease-True[simp]*: $P\ t'.\ t'\ \mathcal{R}\ t\ I.\ \text{True}$
 $\langle\text{proof}\rangle$

lemma *iRelease-False-left[simp]*: $(\text{False}.\ t'\ \mathcal{R}\ t\ I.\ Q\ t) = (\Box\ t\ I.\ Q\ t)$
 $\langle\text{proof}\rangle$

lemma *iRelease-False[simp]*: $(P\ t'.\ t'\ \mathcal{R}\ t\ I.\ \text{False}) = (I = \{\})$
 $\langle\text{proof}\rangle$

lemma *iTrigger-True-left[simp]*: $(\text{True}.\ t'\ \mathcal{T}\ t\ I.\ Q\ t) = (I = \{\}) \vee (\Diamond\ t1\ I.\ \Box\ t2$
 $(I\ \Downarrow\ t1).\ Q\ t2))$
 $\langle\text{proof}\rangle$

lemma *iTrigger-True[simp]*: $P\ t'.\ t'\ \mathcal{T}\ t\ I.\ \text{True}$
 $\langle\text{proof}\rangle$

lemma *iTrigger-False-left[simp]*: $(\text{False}.\ t'\ \mathcal{T}\ t\ I.\ Q\ t) = (\Box\ t\ I.\ Q\ t)$
 $\langle\text{proof}\rangle$

lemma *iTrigger-False[simp]*: $(P\ t'.\ t'\ \mathcal{T}\ t\ I.\ \text{False}) = (I = \{\})$
 $\langle\text{proof}\rangle$

lemma
iUntil-TrueTrue[simp]: $(\text{True}.\ t'\ \mathcal{U}\ t\ I.\ \text{True}) = (I \neq \{\})$ **and**

iSince-TrueTrue[simp]: $(\text{True}. t' \mathcal{S} t I. \text{True}) = (I \neq \{\})$ **and**
iWeakUntil-TrueTrue[simp]: $\text{True}. t' \mathcal{W} t I. \text{True}$ **and**
iWeakSince-TrueTrue[simp]: $\text{True}. t' \mathcal{B} t I. \text{True}$ **and**
iRelease-TrueTrue[simp]: $\text{True}. t' \mathcal{R} t I. \text{True}$ **and**
iTrigger-TrueTrue[simp]: $\text{True}. t' \mathcal{T} t I. \text{True}$
 ⟨proof⟩

3.2.3 Empty sets and singletons

lemma *iAll-empty[simp]*: $\square t \{\}$. $P t$ ⟨proof⟩

lemma *iEx-empty[simp]*: $\neg (\diamond t \{\})$. $P t$ ⟨proof⟩

lemma *iUntil-empty[simp]*: $\neg (P t0. t0 \mathcal{U} t1 \{\}). Q t1$ ⟨proof⟩

lemma *iSince-empty[simp]*: $\neg (P t0. t0 \mathcal{S} t1 \{\}). Q t1$ ⟨proof⟩

lemma *iWeakUntil-empty[simp]*: $P t0. t0 \mathcal{W} t1 \{\}). Q t1$ ⟨proof⟩

lemma *iWeakSince-empty[simp]*: $P t0. t0 \mathcal{B} t1 \{\}). Q t1$ ⟨proof⟩

lemma *iRelease-empty[simp]*: $P t0. t0 \mathcal{R} t1 \{\}). Q t1$ ⟨proof⟩

lemma *iTrigger-empty[simp]*: $P t0. t0 \mathcal{T} t1 \{\}). Q t1$ ⟨proof⟩

lemmas *iTL-empty* =

iAll-empty iEx-empty

iUntil-empty iSince-empty

iWeakUntil-empty iWeakSince-empty

iRelease-empty iTrigger-empty

lemma *iAll-singleton[simp]*: $(\square t' \{t\}. P t') = P t$ ⟨proof⟩

lemma *iEx-singleton[simp]*: $(\diamond t' \{t\}. P t') = P t$ ⟨proof⟩

lemma *iUntil-singleton[simp]*: $(P t0. t0 \mathcal{U} t1 \{t\}. Q t1) = Q t$
 ⟨proof⟩

lemma *iSince-singleton[simp]*: $(P t0. t0 \mathcal{S} t1 \{t\}. Q t1) = Q t$
 ⟨proof⟩

lemma *iWeakUntil-singleton[simp]*: $(P t0. t0 \mathcal{W} t1 \{t\}. Q t1) = (P t \vee Q t)$
 ⟨proof⟩

lemma *iWeakSince-singleton[simp]*: $(P t0. t0 \mathcal{B} t1 \{t\}. Q t1) = (P t \vee Q t)$
 ⟨proof⟩

lemma *iRelease-singleton[simp]*: $(P t0. t0 \mathcal{R} t1 \{t\}. Q t1) = Q t$
 ⟨proof⟩

lemma *iTrigger-singleton[simp]*: $(P t0. t0 \mathcal{T} t1 \{t\}. Q t1) = Q t$
 ⟨proof⟩

lemmas *iTL-singleton* =

iAll-singleton iEx-singleton

iUntil-singleton iSince-singleton
iWeakUntil-singleton iWeakSince-singleton
iRelease-singleton iTrigger-singleton

3.2.4 Conversions between temporal operators

lemma *iAll-iEx-conv*: $(\Box t I. P t) = (\neg (\Diamond t I. \neg P t))$ *<proof>*

lemma *iEx-iAll-conv*: $(\Diamond t I. P t) = (\neg (\Box t I. \neg P t))$ *<proof>*

lemma *not-iAll[simp]*: $(\neg (\Box t I. P t)) = (\Diamond t I. \neg P t)$ *<proof>*

lemma *not-iEx[simp]*: $(\neg (\Diamond t I. P t)) = (\Box t I. \neg P t)$ *<proof>*

lemma *iUntil-iEx-conv*: $(\text{True}. t' \mathcal{U} t I. P t) = (\Diamond t I. P t)$ *<proof>*

lemma *iSince-iEx-conv*: $(\text{True}. t' \mathcal{S} t I. P t) = (\Diamond t I. P t)$ *<proof>*

lemma *iRelease-iAll-conv*: $(\text{False}. t' \mathcal{R} t I. P t) = (\Box t I. P t)$
<proof>

lemma *iTrigger-iAll-conv*: $(\text{False}. t' \mathcal{T} t I. P t) = (\Box t I. P t)$
<proof>

lemma *iWeakUntil-iUntil-conv*: $(P t'. t' \mathcal{W} t I. Q t) = ((P t'. t' \mathcal{U} t I. Q t) \vee (\Box t I. P t))$
<proof>

lemma *iWeakSince-iSince-conv*: $(P t'. t' \mathcal{B} t I. Q t) = ((P t'. t' \mathcal{S} t I. Q t) \vee (\Box t I. P t))$
<proof>

lemma *iUntil-iWeakUntil-conv*: $(P t'. t' \mathcal{U} t I. Q t) = ((P t'. t' \mathcal{W} t I. Q t) \wedge (\Diamond t I. Q t))$
<proof>

lemma *iSince-iWeakSince-conv*: $(P t'. t' \mathcal{S} t I. Q t) = ((P t'. t' \mathcal{B} t I. Q t) \wedge (\Diamond t I. Q t))$
<proof>

lemma *iRelease-iWeakUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = (Q t'. t' \mathcal{W} t I. (Q t \wedge P t))$
<proof>

lemma *iRelease-iUntil-conv*: $(P t'. t' \mathcal{R} t I. Q t) = ((\Box t I. Q t) \vee (Q t'. t' \mathcal{U} t I. (Q t \wedge P t)))$
<proof>

lemma *iTrigger-iWeakSince-conv*: $(P t'. t' \mathcal{T} t I. Q t) = (Q t'. t' \mathcal{B} t I. (Q t \wedge P t))$
<proof>

lemma *iTrigger-iSince-conv*: $(P\ t'.\ t'\ \mathcal{T}\ t\ I.\ Q\ t) = ((\Box\ t\ I.\ Q\ t) \vee (Q\ t'.\ t'\ \mathcal{S}\ t\ I.\ (Q\ t \wedge P\ t)))$
 ⟨proof⟩

lemma *iRelease-not-iUntil-conv*: $(P\ t'.\ t'\ \mathcal{R}\ t\ I.\ Q\ t) = (\neg(\neg P\ t'.\ t'\ \mathcal{U}\ t\ I.\ \neg Q\ t))$
 ⟨proof⟩

lemma *iUntil-not-iRelease-conv*: $(P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t) = (\neg(\neg P\ t'.\ t'\ \mathcal{R}\ t\ I.\ \neg Q\ t))$
 ⟨proof⟩

The Trigger operator \mathcal{T} is a past operator, so that it is used for time intervals, that are bounded by a current time point, and thus are finite. For an infinite interval the stated relation to the Since operator \mathcal{S} would not be fulfilled.

lemma *iTrigger-not-iSince-conv*: $\text{finite } I \implies (P\ t'.\ t'\ \mathcal{T}\ t\ I.\ Q\ t) = (\neg(\neg P\ t'.\ t'\ \mathcal{S}\ t\ I.\ \neg Q\ t))$
 ⟨proof⟩

lemma *iSince-not-iTrigger-conv*: $\text{finite } I \implies (P\ t'.\ t'\ \mathcal{S}\ t\ I.\ Q\ t) = (\neg(\neg P\ t'.\ t'\ \mathcal{T}\ t\ I.\ \neg Q\ t))$
 ⟨proof⟩

lemma *not-iUntil*:
 $(\neg(P\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2)) =$
 $(\Box\ t\ I.\ (Q\ t \longrightarrow (\Diamond\ t'\ (I\ \downarrow < t).\ \neg P\ t')))$
 ⟨proof⟩

lemma *not-iSince*:
 $(\neg(P\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2)) =$
 $(\Box\ t\ I.\ (Q\ t \longrightarrow (\Diamond\ t'\ (I\ \downarrow > t).\ \neg P\ t')))$
 ⟨proof⟩

lemma *iWeakUntil-conj-iUntil-conv*:
 $(P\ t1.\ t1\ \mathcal{W}\ t2\ I.\ (P\ t2 \wedge Q\ t2)) = (\neg(\neg Q\ t1.\ t1\ \mathcal{U}\ t2\ I.\ \neg P\ t2))$
 ⟨proof⟩

lemma *iUntil-disj-iUntil-conv*:
 $(P\ t1 \vee Q\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2) =$
 $(P\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2)$
 ⟨proof⟩

lemma *iWeakUntil-disj-iWeakUntil-conv*:
 $(P\ t1 \vee Q\ t1.\ t1\ \mathcal{W}\ t2\ I.\ Q\ t2) =$
 $(P\ t1.\ t1\ \mathcal{W}\ t2\ I.\ Q\ t2)$
 ⟨proof⟩

lemma *iWeakUntil-iUntil-conj-conv*:

$$(P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2) = \\ (\neg (\neg Q \ t1. \ t1 \ \mathcal{U} \ t2 \ I. (\neg P \ t2 \wedge \neg Q \ t2))) \\ \langle \text{proof} \rangle$$

Negation and temporal operators

lemma

$$\begin{aligned} \text{not-iNext[simp]: } & (\neg (\bigcirc \ t \ t0 \ I. \ P \ t)) = (\bigcirc \ t \ t0 \ I. \ \neg \ P \ t) \ \mathbf{and} \\ \text{not-iNextWeak[simp]: } & (\neg (\bigcirc_W \ t \ t0 \ I. \ P \ t)) = (\bigcirc_S \ t \ t0 \ I. \ \neg \ P \ t) \ \mathbf{and} \\ \text{not-iNextStrong[simp]: } & (\neg (\bigcirc_S \ t \ t0 \ I. \ P \ t)) = (\bigcirc_W \ t \ t0 \ I. \ \neg \ P \ t) \ \mathbf{and} \\ \text{not-iLast[simp]: } & (\neg (\ominus \ t \ t0 \ I. \ P \ t)) = (\ominus \ t \ t0 \ I. \ \neg \ P \ t) \ \mathbf{and} \\ \text{not-iLastWeak[simp]: } & (\neg (\ominus_W \ t \ t0 \ I. \ P \ t)) = (\ominus_S \ t \ t0 \ I. \ \neg \ P \ t) \ \mathbf{and} \\ \text{not-iLastStrong[simp]: } & (\neg (\ominus_S \ t \ t0 \ I. \ P \ t)) = (\ominus_W \ t \ t0 \ I. \ \neg \ P \ t) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *not-iWeakUntil*:

$$\begin{aligned} & (\neg (P \ t1. \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ t2)) = \\ & ((\Box \ t \ I. (Q \ t \longrightarrow (\Diamond \ t' (I \ \downarrow < \ t). \ \neg \ P \ t'))) \wedge (\Diamond \ t \ I. \ \neg \ P \ t)) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *not-iWeakSince*:

$$\begin{aligned} & (\neg (P \ t1. \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ t2)) = \\ & ((\Box \ t \ I. (Q \ t \longrightarrow (\Diamond \ t' (I \ \downarrow > \ t). \ \neg \ P \ t'))) \wedge (\Diamond \ t \ I. \ \neg \ P \ t)) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *not-iRelease*:

$$\begin{aligned} & (\neg (P \ t'. \ t' \ \mathcal{R} \ t \ I. \ Q \ t)) = \\ & ((\Diamond \ t \ I. \ \neg \ Q \ t) \wedge (\Box \ t \ I. \ P \ t \longrightarrow (\Diamond \ t \ I \ \downarrow \leq \ t. \ \neg \ Q \ t))) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *not-iRelease-iUntil-conv*:

$$\begin{aligned} & (\neg (P \ t'. \ t' \ \mathcal{R} \ t \ I. \ Q \ t)) = (\neg P \ t'. \ t' \ \mathcal{U} \ t \ I. \ \neg \ Q \ t) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *not-iTrigger*:

$$\begin{aligned} & (\neg (P \ t'. \ t' \ \mathcal{T} \ t \ I. \ Q \ t)) = \\ & ((\Diamond \ t \ I. \ \neg \ Q \ t) \wedge (\Box \ t \ I. \ \neg \ P \ t \vee (\Diamond \ t \ I \ \downarrow \geq \ t. \ \neg \ Q \ t))) \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma *not-iTrigger-iSince-conv*:

$$\begin{aligned} & \text{finite } I \implies (\neg (P \ t'. \ t' \ \mathcal{T} \ t \ I. \ Q \ t)) = (\neg P \ t'. \ t' \ \mathcal{S} \ t \ I. \ \neg \ Q \ t) \\ & \langle \text{proof} \rangle \end{aligned}$$

3.2.5 Some implication results

lemma *all-imp-iall*: $\forall x. \ P \ x \implies \Box \ t \ I. \ P \ t$ *<proof>*

lemma *bex-imp-lex*: $\Diamond \ t \ I. \ P \ t \implies \exists x. \ P \ x$ *<proof>*

lemma *iAll-imp-iEx*: $I \neq \{\} \implies \Box \ t \ I. \ P \ t \implies \Diamond \ t \ I. \ P \ t$ *<proof>*

lemma *i-set-iAll-imp-iEx*: $I \in \text{i-set} \implies \Box \ t \ I. \ P \ t \implies \Diamond \ t \ I. \ P \ t$

$\langle proof \rangle$

lemmas $iT\text{-}iAll\text{-}imp\text{-}iEx = iT\text{-}not\text{-}empty[THEN\ iAll\text{-}imp\text{-}iEx]$

lemma $iUntil\text{-}imp\text{-}iEx: P\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2 \implies \diamond\ t\ I.\ Q\ t$
 $\langle proof \rangle$

lemma $iSince\text{-}imp\text{-}iEx: P\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2 \implies \diamond\ t\ I.\ Q\ t$
 $\langle proof \rangle$

lemma $iAll\text{-}subset\text{-}imp\text{-}iAll: \llbracket \square\ t\ B.\ P\ t; A \subseteq B \rrbracket \implies \square\ t\ A.\ P\ t$
 $\langle proof \rangle$

lemma $iEx\text{-}subset\text{-}imp\text{-}iEx: \llbracket \diamond\ t\ A.\ P\ t; A \subseteq B \rrbracket \implies \diamond\ t\ B.\ P\ t$
 $\langle proof \rangle$

lemma $iAll\text{-}mp: \llbracket \square\ t\ I.\ P\ t \longrightarrow Q\ t; \square\ t\ I.\ P\ t \rrbracket \implies \square\ t\ I.\ Q\ t$ $\langle proof \rangle$

lemma $iEx\text{-}mp: \llbracket \square\ t\ I.\ P\ t \longrightarrow Q\ t; \diamond\ t\ I.\ P\ t \rrbracket \implies \diamond\ t\ I.\ Q\ t$ $\langle proof \rangle$

lemma $iUntil\text{-}imp:$

$\llbracket P1\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2; \square\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \implies P2\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma $iSince\text{-}imp:$

$\llbracket P1\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2; \square\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \implies P2\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma $iWeakUntil\text{-}imp:$

$\llbracket P1\ t1.\ t1\ \mathcal{W}\ t2\ I.\ Q\ t2; \square\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \implies P2\ t1.\ t1\ \mathcal{W}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma $iWeakSince\text{-}imp:$

$\llbracket P1\ t1.\ t1\ \mathcal{B}\ t2\ I.\ Q\ t2; \square\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \implies P2\ t1.\ t1\ \mathcal{B}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma $iRelease\text{-}imp:$

$\llbracket P1\ t1.\ t1\ \mathcal{R}\ t2\ I.\ Q\ t2; \square\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \implies P2\ t1.\ t1\ \mathcal{R}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma $iTrigger\text{-}imp:$

$\llbracket P1\ t1.\ t1\ \mathcal{T}\ t2\ I.\ Q\ t2; \square\ t\ I.\ P1\ t \longrightarrow P2\ t \rrbracket \implies P2\ t1.\ t1\ \mathcal{T}\ t2\ I.\ Q\ t2$
 $\langle proof \rangle$

lemma $iMin\text{-}imp\text{-}iUntil:$

$\llbracket I \neq \{\}; Q\ (iMin\ I) \rrbracket \implies P\ t'.\ t'\ \mathcal{U}\ t\ I.\ Q\ t$
 $\langle proof \rangle$

lemma *Max-imp-iSince*:

$$\llbracket \text{finite } I; I \neq \{\}; Q (\text{Max } I) \rrbracket \implies P t'. t' \mathcal{S} t I. Q t$$

<proof>

3.2.6 Congruence rules for temporal operators' predicates

lemma *iAll-cong*: $\Box t I. f t = g t \implies (\Box t I. P (f t) t) = (\Box t I. P (g t) t)$

<proof>

lemma *iEx-cong*: $\Box t I. f t = g t \implies (\Diamond t I. P (f t) t) = (\Diamond t I. P (g t) t)$

<proof>

lemma *iUntil-cong1*:

$$\Box t I. f t = g t \implies (P (f t1) t1. t1 \mathcal{U} t2 I. Q t2) = (P (g t1) t1. t1 \mathcal{U} t2 I. Q t2)$$

<proof>

lemma *iUntil-cong2*:

$$\Box t I. f t = g t \implies (P t1. t1 \mathcal{U} t2 I. Q (f t2) t2) = (P t1. t1 \mathcal{U} t2 I. Q (g t2) t2)$$

<proof>

lemma *iSince-cong1*:

$$\Box t I. f t = g t \implies (P (f t1) t1. t1 \mathcal{S} t2 I. Q t2) = (P (g t1) t1. t1 \mathcal{S} t2 I. Q t2)$$

<proof>

lemma *iSince-cong2*:

$$\Box t I. f t = g t \implies (P t1. t1 \mathcal{S} t2 I. Q (f t2) t2) = (P t1. t1 \mathcal{S} t2 I. Q (g t2) t2)$$

<proof>

lemma *bex-subst*:

$$\forall x \in A. P x \longrightarrow (Q x = Q' x) \implies (\exists x \in A. P x \wedge Q x) = (\exists x \in A. P x \wedge Q' x)$$

<proof>

lemma *iEx-subst*:

$$\Box t I. P t \longrightarrow (Q t = Q' t) \implies (\Diamond t I. P t \wedge Q t) = (\Diamond t I. P t \wedge Q' t)$$

<proof>

3.2.7 Temporal operators with set unions/intersections and subsets

lemma *iAll-subset*: $\llbracket A \subseteq B; \Box t B. P t \rrbracket \implies \Box t A. P t$

<proof>

lemma *iEx-subset*: $\llbracket A \subseteq B; \Diamond t A. P t \rrbracket \implies \Diamond t B. P t$

$\langle \text{proof} \rangle$

lemma *iUntil-append*:

$\llbracket \text{finite } A; \text{Max } A \leq \text{iMin } B \rrbracket \implies$
 $P \ t1. \ t1 \ \mathcal{U} \ t2 \ A. \ Q \ t2 \implies P \ t1. \ t1 \ \mathcal{U} \ t2 \ (A \cup B). \ Q \ t2$
 $\langle \text{proof} \rangle$

lemma *iSince-prepend*:

$\llbracket \text{finite } A; \text{Max } A \leq \text{iMin } B \rrbracket \implies$
 $P \ t1. \ t1 \ \mathcal{S} \ t2 \ B. \ Q \ t2 \implies P \ t1. \ t1 \ \mathcal{S} \ t2 \ (A \cup B). \ Q \ t2$
 $\langle \text{proof} \rangle$

lemma

iAll-union: $\llbracket \Box \ t \ A. \ P \ t; \Box \ t \ B. \ P \ t \rrbracket \implies \Box \ t \ (A \cup B). \ P \ t$ **and**
iAll-union-conv: $(\Box \ t \ A \cup B. \ P \ t) = ((\Box \ t \ A. \ P \ t) \wedge (\Box \ t \ B. \ P \ t))$
 $\langle \text{proof} \rangle$

lemma

iEx-union: $(\Diamond \ t \ A. \ P \ t) \vee (\Diamond \ t \ B. \ P \ t) \implies \Diamond \ t \ (A \cup B). \ P \ t$ **and**
iEx-union-conv: $(\Diamond \ t \ (A \cup B). \ P \ t) = ((\Diamond \ t \ A. \ P \ t) \vee (\Diamond \ t \ B. \ P \ t))$
 $\langle \text{proof} \rangle$

lemma *iAll-inter*: $(\Box \ t \ A. \ P \ t) \vee (\Box \ t \ B. \ P \ t) \implies \Box \ t \ (A \cap B). \ P \ t$ $\langle \text{proof} \rangle$

lemma *not-iEx-inter*:

$\exists A \ B \ P. (\Diamond \ t \ A. \ P \ t) \wedge (\Diamond \ t \ B. \ P \ t) \wedge \neg (\Diamond \ t \ (A \cap B). \ P \ t)$
 $\langle \text{proof} \rangle$

lemma

iAll-insert: $\llbracket P \ t; \Box \ t \ I. \ P \ t \rrbracket \implies \Box \ t' \ (\text{insert } t \ I). \ P \ t'$ **and**
iAll-insert-conv: $(\Box \ t' \ (\text{insert } t \ I). \ P \ t') = (P \ t \wedge (\Box \ t' \ I. \ P \ t'))$
 $\langle \text{proof} \rangle$

lemma

iEx-insert: $\llbracket P \ t \vee (\Diamond \ t \ I. \ P \ t) \rrbracket \implies \Diamond \ t' \ (\text{insert } t \ I). \ P \ t'$ **and**
iEx-insert-conv: $(\Diamond \ t' \ (\text{insert } t \ I). \ P \ t') = (P \ t \vee (\Diamond \ t' \ I. \ P \ t'))$
 $\langle \text{proof} \rangle$

3.3 Further results for temporal operators

lemma *Collect-minI-iEx*: $\Diamond \ t \ I. \ P \ t \implies \Diamond \ t \ I. \ P \ t \wedge (\Box \ t' \ (I \downarrow < t). \neg P \ t')$
 $\langle \text{proof} \rangle$

lemma *iUntil-disj-conv1*:

$I \neq \{\}$ \implies
 $(P \ t'. \ t' \ \mathcal{U} \ t \ I. \ Q \ t) = (Q \ (iMin \ I) \vee (P \ t'. \ t' \ \mathcal{U} \ t \ I. \ Q \ t \wedge iMin \ I < t))$
 $\langle \text{proof} \rangle$

lemma *iSince-disj-conv1*:

$\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies$
 $(P t'. t' \mathcal{S} t I. Q t) = (Q (\text{Max } I) \vee (P t'. t' \mathcal{S} t I. Q t \wedge t < \text{Max } I))$
 <proof>

lemma *iUntil-next*:

$I \neq \{\} \implies$
 $(P t'. t' \mathcal{U} t I. Q t) =$
 $(Q (\text{iMin } I) \vee (P (\text{iMin } I) \wedge (P t'. t' \mathcal{U} t (I \downarrow > (\text{iMin } I)). Q t)))$
 <proof>

lemma *iSince-prev*: $\llbracket \text{finite } I; I \neq \{\} \rrbracket \implies$

$(P t'. t' \mathcal{S} t I. Q t) =$
 $(Q (\text{Max } I) \vee (P (\text{Max } I) \wedge (P t'. t' \mathcal{S} t (I \downarrow < \text{Max } I). Q t)))$
 <proof>

lemma *iNext-induct-rule*:

$\llbracket P (\text{iMin } I); \square t I. (P t \longrightarrow (\bigcirc t' t I. P t')) \rrbracket \implies P t$
 <proof>

lemma *iNext-induct*:

$\llbracket P (\text{iMin } I); \square t I. (P t \longrightarrow (\bigcirc t' t I. P t')) \rrbracket \implies \square t I. P t$
 <proof>

lemma *iLast-induct-rule*:

$\llbracket P (\text{Max } I); \square t I. (P t \longrightarrow (\ominus t' t I. P t')) \rrbracket; \text{finite } I; t \in I \implies P t$
 <proof>

lemma *iLast-induct*:

$\llbracket P (\text{Max } I); \square t I. (P t \longrightarrow (\ominus t' t I. P t')) \rrbracket; \text{finite } I \implies \square t I. P t$
 <proof>

lemma *iUntil-conj-not*: $((P t1 \wedge \neg Q t1). t1 \mathcal{U} t2 I. Q t2) = (P t1. t1 \mathcal{U} t2 I. Q t2)$

<proof>

lemma *iWeakUntil-conj-not*: $((P t1 \wedge \neg Q t1). t1 \mathcal{W} t2 I. Q t2) = (P t1. t1 \mathcal{W} t2 I. Q t2)$

<proof>

lemma *iSince-conj-not*: $\text{finite } I \implies$

$((P t1 \wedge \neg Q t1). t1 \mathcal{S} t2 I. Q t2) = (P t1. t1 \mathcal{S} t2 I. Q t2)$
 <proof>

lemma *iWeakSince-conj-not*: $\text{finite } I \implies$

$((P t1 \wedge \neg Q t1). t1 \mathcal{B} t2 I. Q t2) = (P t1. t1 \mathcal{B} t2 I. Q t2)$
 <proof>

lemma *iNextStrong-imp-iNextWeak*: $(\circ_S t t0 I. P t) \longrightarrow (\circ_W t t0 I. P t)$
 ⟨proof⟩

lemma *iLastStrong-imp-iLastWeak*: $(\ominus_S t t0 I. P t) \longrightarrow (\ominus_W t t0 I. P t)$
 ⟨proof⟩

lemma *infin-imp-iNextWeak-iNextStrong-eq-iNext*:

[[*infinite I*; $t0 \in I$]] \implies
 $((\circ_W t t0 I. P t) = (\circ t t0 I. P t)) \wedge ((\circ_S t t0 I. P t) = (\circ t t0 I. P t))$
 ⟨proof⟩

lemma *infin-imp-iNextWeak-eq-iNext*: [[*infinite I*; $t0 \in I$]] $\implies (\circ_W t t0 I. P t)$
 $= (\circ t t0 I. P t)$

⟨proof⟩

lemma *infin-imp-iNextStrong-eq-iNext*: [[*infinite I*; $t0 \in I$]] $\implies (\circ_S t t0 I. P t)$
 $= (\circ t t0 I. P t)$

⟨proof⟩

lemma *infin-imp-iNextStrong-eq-iNextWeak*: [[*infinite I*; $t0 \in I$]] $\implies (\circ_S t t0 I.$
 $P t) = (\circ_W t t0 I. P t)$

⟨proof⟩

lemma

not-in-iNext-eq: $t0 \notin I \implies (\circ t t0 I. P t) = (P t0)$ **and**

not-in-iLast-eq: $t0 \notin I \implies (\ominus t t0 I. P t) = (P t0)$

⟨proof⟩

lemma

not-in-iNextWeak-eq: $t0 \notin I \implies (\circ_W t t0 I. P t)$ **and**

not-in-iLastWeak-eq: $t0 \notin I \implies (\ominus_W t t0 I. P t)$

⟨proof⟩

lemma

not-in-iNextStrong-eq: $t0 \notin I \implies \neg (\circ_S t t0 I. P t)$ **and**

not-in-iLastStrong-eq: $t0 \notin I \implies \neg (\ominus_S t t0 I. P t)$

⟨proof⟩

lemma

iNext-UNIV: $(\circ t t0 UNIV. P t) = P (Suc t0)$ **and**

iNextWeak-UNIV: $(\circ_W t t0 UNIV. P t) = P (Suc t0)$ **and**

iNextStrong-UNIV: $(\circ_S t t0 UNIV. P t) = P (Suc t0)$

⟨proof⟩

lemma

iLast-UNIV: $(\ominus t t0 UNIV. P t) = P (t0 - Suc 0)$ **and**

iLastWeak-UNIV: $(\ominus_W t t0 UNIV. P t) = (if\ 0 < t0\ then\ P\ (t0 - Suc\ 0)\ else\ True)$ **and**

iLastStrong-UNIV: $(\ominus_S t t0 UNIV. P t) = (if\ 0 < t0\ then\ P\ (t0 - Suc\ 0)\ else\ False)$

⟨proof⟩

lemmas *iTL-Next-UNIV* =
iNext-UNIV iNextWeak-UNIV iNextStrong-UNIV
iLast-UNIV iLastWeak-UNIV iLastStrong-UNIV

lemma *inext-nth-iNext-Suc*: $(\circ t (I \rightarrow n) I. P t) = P (I \rightarrow \text{Suc } n)$
 $\langle \text{proof} \rangle$

lemma *iprev-nth-iLast-Suc*: $(\ominus t (I \leftarrow n) I. P t) = P (I \leftarrow \text{Suc } n)$
 $\langle \text{proof} \rangle$

3.4 Temporal operators and arithmetic interval operators

Shifting intervals through addition and subtraction of constants. Mirroring intervals through subtraction of intervals from constants. Expanding and compressing intervals through multiplication and division by constants.

Always operator

lemma *iT-Plus-iAll-conv*: $(\Box t I \oplus k. P t) = (\Box t I. P (t + k))$
 $\langle \text{proof} \rangle$

lemma *iT-Mult-iAll-conv*: $(\Box t I \otimes k. P t) = (\Box t I. P (t * k))$
 $\langle \text{proof} \rangle$

lemma *iT-Plus-neg-iAll-conv*: $(\Box t I \oplus - k. P t) = (\Box t (I \downarrow \geq k). P (t - k))$
 $\langle \text{proof} \rangle$

lemma *iT-Minus-iAll-conv*: $(\Box t k \ominus I. P t) = (\Box t (I \downarrow \leq k). P (k - t))$
 $\langle \text{proof} \rangle$

lemma *iT-Div-iAll-conv*: $(\Box t I \oslash k. P t) = (\Box t I. P (t \text{ div } k))$
 $\langle \text{proof} \rangle$

lemmas *iT-arith-iAll-conv* =
iT-Plus-iAll-conv
iT-Mult-iAll-conv
iT-Plus-neg-iAll-conv
iT-Minus-iAll-conv
iT-Div-iAll-conv

Eventually operator

lemma
iT-Plus-iEx-conv: $(\Diamond t I \oplus k. P t) = (\Diamond t I. P (t + k))$ **and**
iT-Mult-iEx-conv: $(\Diamond t I \otimes k. P t) = (\Diamond t I. P (t * k))$ **and**
iT-Plus-neg-iEx-conv: $(\Diamond t I \oplus - k. P t) = (\Diamond t (I \downarrow \geq k). P (t - k))$ **and**
iT-Minus-iEx-conv: $(\Diamond t k \ominus I. P t) = (\Diamond t (I \downarrow \leq k). P (k - t))$ **and**
iT-Div-iEx-conv: $(\Diamond t I \oslash k. P t) = (\Diamond t I. P (t \text{ div } k))$
 $\langle \text{proof} \rangle$

Until and Since operators

lemma *iT-Plus-iUntil-conv*: $(P\ t1.\ t1\ \mathcal{U}\ t2\ (I\ \oplus\ k).\ Q\ t2) = (P\ (t1\ +\ k).\ t1\ \mathcal{U}\ t2\ I.\ Q\ (t2\ +\ k))$
 ⟨proof⟩

lemma *iT-Mult-iUntil-conv*: $(P\ t1.\ t1\ \mathcal{U}\ t2\ (I\ \otimes\ k).\ Q\ t2) = (P\ (t1\ *\ k).\ t1\ \mathcal{U}\ t2\ I.\ Q\ (t2\ *\ k))$
 ⟨proof⟩

lemma *iT-Plus-neg-iUntil-conv*: $(P\ t1.\ t1\ \mathcal{U}\ t2\ (I\ \oplus\ -\ k).\ Q\ t2) = (P\ (t1\ -\ k).\ t1\ \mathcal{U}\ t2\ (I\ \downarrow\geq\ k).\ Q\ (t2\ -\ k))$
 ⟨proof⟩

lemma *iT-Minus-iUntil-conv*: $(P\ t1.\ t1\ \mathcal{U}\ t2\ (k\ \ominus\ I).\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ \mathcal{S}\ t2\ (I\ \downarrow\leq\ k).\ Q\ (k\ -\ t2))$
 ⟨proof⟩

lemma *iT-Div-iUntil-conv*: $(P\ t1.\ t1\ \mathcal{U}\ t2\ (I\ \odot\ k).\ Q\ t2) = (P\ (t1\ \text{div}\ k).\ t1\ \mathcal{U}\ t2\ I.\ Q\ (t2\ \text{div}\ k))$
 ⟨proof⟩

Until and Since operators can be converted into each other through subtraction of intervals from constants

lemma *iUntil-iSince-conv*:
 $\llbracket\ \text{finite}\ I;\ \text{Max}\ I\ \leq\ k\ \rrbracket\ \Longrightarrow$
 $(P\ t1.\ t1\ \mathcal{U}\ t2\ I.\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ \mathcal{S}\ t2\ (k\ \ominus\ I).\ Q\ (k\ -\ t2))$
 ⟨proof⟩

lemma *iSince-iUntil-conv*:
 $\llbracket\ \text{finite}\ I;\ \text{Max}\ I\ \leq\ k\ \rrbracket\ \Longrightarrow$
 $(P\ t1.\ t1\ \mathcal{S}\ t2\ I.\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ \mathcal{U}\ t2\ (k\ \ominus\ I).\ Q\ (k\ -\ t2))$
 ⟨proof⟩

lemma *iT-Plus-iSince-conv*: $(P\ t1.\ t1\ \mathcal{S}\ t2\ (I\ \oplus\ k).\ Q\ t2) = (P\ (t1\ +\ k).\ t1\ \mathcal{S}\ t2\ I.\ Q\ (t2\ +\ k))$
 ⟨proof⟩

lemma *iT-Mult-iSince-conv*: $0 < k \Longrightarrow (P\ t1.\ t1\ \mathcal{S}\ t2\ (I\ \otimes\ k).\ Q\ t2) = (P\ (t1\ *\ k).\ t1\ \mathcal{S}\ t2\ I.\ Q\ (t2\ *\ k))$
 ⟨proof⟩

lemma *iT-Plus-neg-iSince-conv*: $(P\ t1.\ t1\ \mathcal{S}\ t2\ (I\ \oplus\ -\ k).\ Q\ t2) = (P\ (t1\ -\ k).\ t1\ \mathcal{S}\ t2\ (I\ \downarrow\geq\ k).\ Q\ (t2\ -\ k))$
 ⟨proof⟩

lemma *iT-Minus-iSince-conv*:
 $(P\ t1.\ t1\ \mathcal{S}\ t2\ (k\ \ominus\ I).\ Q\ t2) = (P\ (k\ -\ t1).\ t1\ \mathcal{U}\ t2\ (I\ \downarrow\leq\ k).\ Q\ (k\ -\ t2))$
 ⟨proof⟩

lemma *iT-Div-iSince-conv*:

$0 < k \implies (P \ t1. \ t1 \ \mathcal{S} \ t2 \ (I \circlearrowleft k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{S} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 ⟨proof⟩

Weak Until and Weak Since operators

lemma *iT-Plus-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ (t2 + k))$
 ⟨proof⟩

lemma *iT-Mult-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ (t2 * k))$
 ⟨proof⟩

lemma *iT-Plus-neg-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{W} \ t2 \ (I \ \downarrow \geq k). \ Q \ (t2 - k))$
 ⟨proof⟩

lemma *iT-Minus-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{B} \ t2 \ (I \ \downarrow \leq k). \ Q \ (k - t2))$
 ⟨proof⟩

lemma *iT-Div-iWeakUntil-conv*: $(P \ t1. \ t1 \ \mathcal{W} \ t2 \ (I \circlearrowleft k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{W} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 ⟨proof⟩

lemma *iT-Plus-iWeakSince-conv*: $(P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ (t2 + k))$
 ⟨proof⟩

lemma *iT-Mult-iWeakSince-conv*: $0 < k \implies (P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ (t2 * k))$
 ⟨proof⟩

lemma *iT-Plus-neg-iWeakSince-conv*: $(P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{B} \ t2 \ (I \ \downarrow \geq k). \ Q \ (t2 - k))$
 ⟨proof⟩

lemma *iT-Minus-iWeakSince-conv*:
 $(P \ t1. \ t1 \ \mathcal{B} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{W} \ t2 \ (I \ \downarrow \leq k). \ Q \ (k - t2))$
 ⟨proof⟩

lemma *iT-Div-iWeakSince-conv*:
 $0 < k \implies (P \ t1. \ t1 \ \mathcal{B} \ t2 \ (I \circlearrowleft k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{B} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 ⟨proof⟩

Release and Trigger operators

lemma *iT-Plus-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1$

$\mathcal{R} \ t2 \ I. \ Q \ (t2 + k)$
 $\langle proof \rangle$

lemma *iT-Mult-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{R} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{R} \ t2 \ (I \downarrow \geq k). \ Q \ (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{T} \ t2 \ (I \downarrow \leq k). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iRelease-conv*: $(P \ t1. \ t1 \ \mathcal{R} \ t2 \ (I \circ k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{R} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 $\langle proof \rangle$

lemma *iT-Plus-iTrigger-conv*: $(P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \oplus k). \ Q \ t2) = (P \ (t1 + k). \ t1 \ \mathcal{T} \ t2 \ I. \ Q \ (t2 + k))$
 $\langle proof \rangle$

lemma *iT-Mult-iTrigger-conv*: $0 < k \implies (P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \otimes k). \ Q \ t2) = (P \ (t1 * k). \ t1 \ \mathcal{T} \ t2 \ I. \ Q \ (t2 * k))$
 $\langle proof \rangle$

lemma *iT-Plus-neg-iTrigger-conv*: $(P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \oplus - k). \ Q \ t2) = (P \ (t1 - k). \ t1 \ \mathcal{T} \ t2 \ (I \downarrow \geq k). \ Q \ (t2 - k))$
 $\langle proof \rangle$

lemma *iT-Minus-iTrigger-conv*:
 $(P \ t1. \ t1 \ \mathcal{T} \ t2 \ (k \ominus I). \ Q \ t2) = (P \ (k - t1). \ t1 \ \mathcal{R} \ t2 \ (I \downarrow \leq k). \ Q \ (k - t2))$
 $\langle proof \rangle$

lemma *iT-Div-iTrigger-conv*:
 $0 < k \implies (P \ t1. \ t1 \ \mathcal{T} \ t2 \ (I \circ k). \ Q \ t2) = (P \ (t1 \ \text{div} \ k). \ t1 \ \mathcal{T} \ t2 \ I. \ Q \ (t2 \ \text{div} \ k))$
 $\langle proof \rangle$

end