

The Nash-Williams Theorem

Lawrence C. Paulson

February 23, 2021

Abstract

In 1965, Nash-Williams [1] discovered a generalisation of the infinite form of Ramsey’s theorem. Where the latter concerns infinite sets of n -element sets for some fixed n , the Nash-Williams theorem concerns infinite sets of finite sets (or lists) subject to a “no initial segment” condition. The present formalisation follows Todorčević [2].

Contents

1	The Pointwise Less-Than Relation Between Two Sets	1
2	The Nash-Williams Theorem	2
2.1	Initial segments	3
2.2	Definitions and basic properties	3
2.3	Main Theorem	6
3	Acknowledgements	6

1 The Pointwise Less-Than Relation Between Two Sets

theory *Nash-Extras*

imports *HOL-Library.Ramsey HOL-Library.Countable-Set*

begin

definition *less-sets* :: [*'a::order set, 'a::order set*] \Rightarrow *bool* **where**
less-sets *A B* $\equiv \forall x \in A. \forall y \in B. x < y$

lemma *less-setsD*: $[[\textit{less-sets } A B; a \in A; b \in B]] \Longrightarrow a < b$
<proof>

lemma *less-sets-irrefl* [*simp*]: $\textit{less-sets } A A \longleftrightarrow A = \{\}$
<proof>

lemma *less-sets-trans*: $\llbracket \text{less-sets } A B; \text{less-sets } B C; B \neq \{\} \rrbracket \implies \text{less-sets } A C$
<proof>

lemma *less-sets-weaken1*: $\llbracket \text{less-sets } A' B; A \subseteq A' \rrbracket \implies \text{less-sets } A B$
<proof>

lemma *less-sets-weaken2*: $\llbracket \text{less-sets } A B'; B \subseteq B' \rrbracket \implies \text{less-sets } A B$
<proof>

lemma *less-sets-imp-disjnt*: $\text{less-sets } A B \implies \text{disjnt } A B$
<proof>

lemma *less-sets-UN1*: $\text{less-sets } (\bigcup \mathcal{A}) B \longleftrightarrow (\forall A \in \mathcal{A}. \text{less-sets } A B)$
<proof>

lemma *less-sets-UN2*: $\text{less-sets } A (\bigcup \mathcal{B}) \longleftrightarrow (\forall B \in \mathcal{B}. \text{less-sets } A B)$
<proof>

lemma *less-sets-Un1*: $\text{less-sets } (A \cup A') B \longleftrightarrow \text{less-sets } A B \wedge \text{less-sets } A' B$
<proof>

lemma *less-sets-Un2*: $\text{less-sets } A (B \cup B') \longleftrightarrow \text{less-sets } A B \wedge \text{less-sets } A B'$
<proof>

lemma *strict-sorted-imp-less-sets*:
 $\text{strict-sorted } (as @ bs) \implies \text{less-sets } (\text{list.set } as) (\text{list.set } bs)$
<proof>

lemma *Sup-nat-less-sets-singleton*:
fixes $n :: \text{nat}$
assumes $\text{Sup } T < n$ *finite* T
shows $\text{less-sets } T \{n\}$
<proof>

end

2 The Nash-Williams Theorem

Following S. Todorćević, *Introduction to Ramsey Spaces*, Princeton University Press (2010), 11–12.

theory *Nash-Williams*
imports *Nash-Extras*
begin

lemma *finite-nat-Int-greaterThan-iff*:
fixes $N :: \text{nat set}$
shows $\text{finite } (N \cap \{n < ..\}) \longleftrightarrow \text{finite } N$
<proof>

2.1 Initial segments

definition *init-segment* :: *nat set* \Rightarrow *nat set* \Rightarrow *bool*
where *init-segment* *S T* $\equiv \exists S'. T = S \cup S' \wedge \text{less-sets } S S'$

lemma *init-segment-subset*: *init-segment* *S T* $\Longrightarrow S \subseteq T$
 $\langle \text{proof} \rangle$

lemma *init-segment-refl*: *init-segment* *S S*
 $\langle \text{proof} \rangle$

lemma *init-segment-antisym*: $\llbracket \text{init-segment } S T; \text{init-segment } T S \rrbracket \Longrightarrow S=T$
 $\langle \text{proof} \rangle$

lemma *init-segment-trans*: $\llbracket \text{init-segment } S T; \text{init-segment } T U \rrbracket \Longrightarrow \text{init-segment } S U$
 $\langle \text{proof} \rangle$

lemma *init-segment-empty2* [*iff*]: *init-segment* *S* $\{\}$ $\longleftrightarrow S=\{\}$
 $\langle \text{proof} \rangle$

lemma *init-segment-Un*: *less-sets* *S S'* $\Longrightarrow \text{init-segment } S (S \cup S')$
 $\langle \text{proof} \rangle$

lemma *init-segment-iff*:
shows *init-segment* *S T* $\longleftrightarrow S=T \vee (\exists m \in T. S = \{n \in T. n < m\})$ (**is**
 $?lhs=?rhs$)
 $\langle \text{proof} \rangle$

lemma *init-segment-empty* [*iff*]: *init-segment* $\{\}$ *S*
 $\langle \text{proof} \rangle$

lemma *init-segment-insert-iff*:
assumes *Sn*: *less-sets* *S* $\{n\}$ **and** *TS*: $\bigwedge x. x \in T-S \Longrightarrow n \leq x$
shows *init-segment* (*insert* *n S*) *T* $\longleftrightarrow \text{init-segment } S T \wedge n \in T$
 $\langle \text{proof} \rangle$

lemma *init-segment-insert*:
assumes *init-segment* *S T* **and** *T*: *less-sets* *T* $\{n\}$
shows *init-segment* *S* (*insert* *n T*)
 $\langle \text{proof} \rangle$

2.2 Definitions and basic properties

definition *Ramsey* :: [*nat set set*, *nat*] \Rightarrow *bool*
where *Ramsey* $\mathcal{F} r \equiv \forall f \in \mathcal{F} \rightarrow \{..<r\}.$
 $\forall M. \text{infinite } M \longrightarrow$
 $(\exists N i. N \subseteq M \wedge \text{infinite } N \wedge i < r \wedge (\forall j < r. j \neq i \longrightarrow f - \{j\} \cap \mathcal{F} \cap \text{Pow } N = \{\}))$

definition *thin-set* :: nat set set \Rightarrow bool

where *thin-set* $\mathcal{F} \equiv \mathcal{F} \subseteq \text{Collect finite} \wedge (\forall S \in \mathcal{F}. \forall T \in \mathcal{F}. \text{init-segment } S \ T \longrightarrow S=T)$

definition *comparables* :: nat set \Rightarrow nat set \Rightarrow nat set set

where *comparables* $S \ M \equiv \{T. \text{finite } T \wedge (\text{init-segment } T \ S \vee \text{init-segment } S \ T \wedge T-S \subseteq M)\}$

lemma *comparables-iff*: $T \in \text{comparables } S \ M \longleftrightarrow \text{finite } T \wedge (\text{init-segment } T \ S \vee \text{init-segment } S \ T \wedge T \subseteq S \cup M)$

<proof>

lemma *comparables-subset*: $\bigcup (\text{comparables } S \ M) \subseteq S \cup M$

<proof>

lemma *comparables-empty* [*simp*]: $\text{comparables } \{\} \ M = \text{Fpow } M$

<proof>

lemma *comparables-mono*: $N \subseteq M \implies \text{comparables } S \ N \subseteq \text{comparables } S \ M$

<proof>

definition *rejects* $\mathcal{F} \ S \ M \equiv \text{comparables } S \ M \cap \mathcal{F} = \{\}$

abbreviation *accepts*

where *accepts* $\mathcal{F} \ S \ M \equiv \neg \text{rejects } \mathcal{F} \ S \ M$

definition *strongly-accepts*

where *strongly-accepts* $\mathcal{F} \ S \ M \equiv (\forall N \subseteq M. \text{rejects } \mathcal{F} \ S \ N \longrightarrow \text{finite } N)$

definition *decides*

where *decides* $\mathcal{F} \ S \ M \equiv \text{rejects } \mathcal{F} \ S \ M \vee \text{strongly-accepts } \mathcal{F} \ S \ M$

definition *decides-subsets*

where *decides-subsets* $\mathcal{F} \ M \equiv \forall T. T \subseteq M \longrightarrow \text{finite } T \longrightarrow \text{decides } \mathcal{F} \ T \ M$

lemma *strongly-accepts-imp-accepts*:

$\llbracket \text{strongly-accepts } \mathcal{F} \ S \ M; \text{infinite } M \rrbracket \implies \text{accepts } \mathcal{F} \ S \ M$

<proof>

lemma *rejects-trivial*: $\llbracket \text{rejects } \mathcal{F} \ S \ M; \text{thin-set } \mathcal{F}; \text{init-segment } F \ S; F \in \mathcal{F} \rrbracket \implies$

False

<proof>

lemma *rejects-subset*: $\llbracket \text{rejects } \mathcal{F} \ S \ M; N \subseteq M \rrbracket \implies \text{rejects } \mathcal{F} \ S \ N$

<proof>

lemma *strongly-accepts-subset*: $\llbracket \text{strongly-accepts } \mathcal{F} \ S \ M; N \subseteq M \rrbracket \implies \text{strongly-accepts } \mathcal{F} \ S \ N$

<proof>

lemma *decides-subset*: $\llbracket \text{decides } \mathcal{F} S M; N \subseteq M \rrbracket \implies \text{decides } \mathcal{F} S N$
<proof>

lemma *decides-subsets-subset*: $\llbracket \text{decides-subsets } \mathcal{F} M; N \subseteq M \rrbracket \implies \text{decides-subsets } \mathcal{F} N$
<proof>

lemma *rejects-empty [simp]*: $\text{rejects } \mathcal{F} \{ \} M \longleftrightarrow \text{Fpow } M \cap \mathcal{F} = \{ \}$
<proof>

lemma *strongly-accepts-empty [simp]*: $\text{strongly-accepts } \mathcal{F} \{ \} M \longleftrightarrow (\forall N \subseteq M. \text{Fpow } N \cap \mathcal{F} = \{ \} \longrightarrow \text{finite } N)$
<proof>

lemma *ex-infinite-decides-1*:
assumes *infinite M*
obtains *N where N ⊆ M infinite N decides F S N*
<proof>

proposition *ex-infinite-decides-finite*:
assumes *infinite M finite S*
obtains *N where N ⊆ M infinite N ∧ T. T ⊆ S ⇒ decides F T N*
<proof>

lemma *sorted-wrt-subset*: $\llbracket X \in \text{list.set } l; \text{sorted-wrt } (\leq) l \rrbracket \implies \text{hd } l \subseteq X$
<proof>

Todorčević's Lemma 1.18

proposition *ex-infinite-decides-subsets*:
assumes *thin-set F infinite M*
obtains *N where N ⊆ M infinite N decides-subsets F N*
<proof>

Todorčević's Lemma 1.19

proposition *strongly-accepts-1-19*:
assumes *acc: strongly-accepts F S M*
and *thin-set F infinite M S ⊆ M finite S*
and *dsM: decides-subsets F M*
shows *finite {n ∈ M. ¬ strongly-accepts F (insert n S) M}*
<proof>

Much work is needed for this slight strengthening of the previous result!

proposition *strongly-accepts-1-19-plus*:
assumes *thin-set F infinite M*
and *dsM: decides-subsets F M*
obtains *N where N ⊆ M infinite N*

$$\bigwedge S n. \llbracket S \subseteq N; \text{finite } S; \text{strongly-accepts } \mathcal{F} S N; n \in N; \text{less-sets } S \{n\} \rrbracket \\ \implies \text{strongly-accepts } \mathcal{F} (\text{insert } n S) N$$

<proof>

2.3 Main Theorem

Weirdly, the assumption $f \text{ ' } \mathcal{F} \subseteq \{..<2::'a\}$ is not used here; it's perhaps unnecessary due to the particular way that *Ramsey* is defined. It's only needed for $(2::'a) < r$

theorem *Nash-Williams-2:*

assumes *thin-set* \mathcal{F} **shows** *Ramsey* $\mathcal{F} 2$

<proof>

theorem *Nash-Williams:*

assumes \mathcal{F} : *thin-set* $\mathcal{F} r > 0$ **shows** *Ramsey* $\mathcal{F} r$

<proof>

end

3 Acknowledgements

The author was supported by the ERC Advanced Grant ALEXANDRIA (Project 742178) funded by the European Research Council. Todorčević provided help with the proofs by email.

References

- [1] C. S. J. A. Nash-Williams. On well-quasi-ordering transfinite sequences. *Mathematical Proceedings of the Cambridge Philosophical Society*, 61(1):33–39, 1965.
- [2] S. Todorčević. *Introduction to Ramsey Spaces*. Princeton University Press, 2010.