

Morley's Theorem*

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Contents

1	Introduction	2
2	Angles between three complex	2
3	Complex triangles	4
4	Complex trigonometry	5
4.1	The law of sines	7
5	Third unity root	12
6	Axial symmetry in complex field	28
7	Rotations	32
8	Morley's theorem	37

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theory Complex-Angles

imports Complex-Geometry.Elementary-Complex-Geometry

begin

```

1 Introduction

Morley's trisector theorem states that in any triangle, the three points of intersection of the adjacent angle trisectors form an equilateral triangle, called the first Morley triangle or simply the Morley triangle. This theorem is listed in the 100 theorem list [3].

In this theory, we define some basics elements on complex geometry such as axial symmetry, rotations. We also define basic property of congruent triangles in the complex field following the model already presented in the afp [1] In addition we demonstrate sines law in the complex context.

Finally we prove the Morley theorem using Alain Connes's proof [2].

2 Angles between three complex

```
definition angle-c-def:⟨angle-c a b c ≡ ∠(a-b)(c-b)⟩
```

```
lemma angle-c-commute-pi:
assumes ⟨angle-c a b c = pi⟩
shows angle-c a b c = angle-c c b a
using angle-c-def assms ang-vec-sym-pi by presburger
```

```
lemma angle-c-commute:
assumes ⟨angle-c a b c ≠ pi⟩
shows angle-c a b c = -angle-c c b a
using angle-c-def assms ang-vec-sym by presburger
```

```
lemma angle-c-sum:
assumes 1:⟨a - b ≠ 0⟩ and 2:⟨c - b ≠ 0⟩ and 3:⟨b - d ≠ 0⟩
shows ⟨|angle-c a b c + angle-c c b d| = angle-c a b d⟩
using angle-c-def canon-ang-sum by force
```

```
lemma collinear-angle:⟨collinear a b c ⟹ a ≠ b ⟹ b ≠ c ⟹ c ≠ a ⟹ angle-c a b c = 0 ∨ angle-c a b c = pi⟩
unfolding collinear-def unfolding angle-c-def
by (metis ang-vec-def arg-div collinear-def collinear-sym2' is-real-arg2 right-minus-eq)
```

```
lemma cdist-commute:⟨cdist a b = cdist b a⟩
by (simp add: norm-minus-commute)
```

```

lemma angle-sum-triangle:
  assumes h:a ≠ b ∧ b ≠ c ∧ a ≠ c
  shows |∠(c-a)(b-a) + ∠(a-b)(c-b) + ∠(b-c)(a-c)| = pi
proof -
  have f0:<b-a ≠ 0> <c-a≠0> <b-c≠0> using h by auto
  then have ∠c(c-a)(b-a) = abs(Arg((b-a)/(c-a)))
  unfolding ang-vec-def ang-vec-c-def
  apply(subst arg-div[OF f0(1) f0(2), symmetric])
  by(simp)
  have <(b-a)/(c-a) = rcis(cmod((b-a)/(c-a))) (Arg((b-a)/(c-a)))>
  by (simp add: rcis-cmod-Arg)
  have <(b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c) = -1>
  using h
  by (smt (z3) divide-eq-0-iff divide-eq-minus-1-iff f0(1,2,3) minus-diff-eq mult-minus-right
    nonzero-eq-divide-eq times-divide-eq-left)
  then have <Arg((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) = pi>
  by (metis arg-cis cis-pi pi-canonical)
  then have <Arg((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) =
  = |Arg((b-a)/(c-a) * (c-b)/(a-b)) + Arg((a-c)/(b-c))|>
  using arg-mult f0
  by (metis (no-types, lifting) Arg-zero arg-mult divide-divide-eq-left' pi-neq-zero
    times-divide-eq-left
    times-divide-eq-right)
  then have <|Arg((b-a)/(c-a) * (c-b)/(a-b)) + Arg((a-c)/(b-c))| =
  = |Arg((b-a)/(c-a)) + Arg((c-b)/(a-b)) + Arg((a-c)/(b-c))|>
  by (metis (no-types, lifting) arg-mult canon-ang-arg canon-ang-sum divide-eq-0-iff
    eq-iff-diff-eq-0 h
    no-zero-divisors times-divide-eq-right)
  then show ?thesis
  using
    <Arg((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) = |Arg((b-a)/(c-a) * (c-b)/(a-b)) + Arg((a-c)/(b-c))|>
    <Arg((b-a)/(c-a) * (c-b)/(a-b) * (a-c)/(b-c)) = pi> arg-div
  h by force
qed

```

```

lemma angle-sum-triangle-c:
  assumes h:a ≠ b ∧ b ≠ c ∧ a ≠ c
  shows |angle-c c a b + angle-c a b c + angle-c b c a| = pi
  unfolding angle-c-def
  using h by (smt (z3) angle-sum-triangle h)

```

```

lemma angle-sum-triangle':
  assumes h:a ≠ 0 ∧ b ≠ 0 ∧ c ≠ 0
  shows |∠(-a)b + ∠(-b)c + ∠(-c)a| = pi
proof -
  have f0:<a ≠ 0> <b≠0> <c≠0> using h by auto
  have <b/(-a) * a/(-c) * c/(-b) = -1>

```

```

using h by(auto)
then have f10:⟨Arg (b/(-a) * a/(-c) * c/(-b)) = pi⟩
  by (metis arg-cis cis-pi pi-canonical)
then have f1:⟨Arg (b/(-a) * a/(-c) * c/(-b))
= |Arg (b/(-a) * a/(-c)) + Arg (c/(-b))|⟩
  using arg-mult f0
  by (metis ⟨b / - a * a / - c * c / - b = - 1⟩ divide-eq-0-iff divide-eq-minus-1-iff
    times-divide-eq-right)
then have f2:⟨|Arg (b/(-a) * a/(-c)) + Arg (c/(-b))|
= |Arg (b/(-a)) + Arg (a/(-c)) + Arg (c/(-b))|⟩
  by (metis (no-types, lifting) arg-mult canon-ang-arg canon-ang-sum divide-eq-0-iff
    f0(1,2,3)
    neg-0-equal-iff-equal no-zero-divisors times-divide-eq-right)
then show ?thesis
  using f1 f2 arg-div h
  by (smt (verit) f10 add.left-commute ang-vec-def neg-equal-0-iff-equal)
qed

lemma ang-c-in:⟨angle-c a b c ∈ {−pi..pi}⟩
  unfolding angle-c-def by(auto simp: canon-ang)

lemma angle-c-aac:⟨angle-c a a c = |Arg (c-a)|⟩
  by (simp add: Arg-zero angle-c-def)

lemma angle-c-caa:⟨angle-c c a a = |−Arg (c-a)|⟩
  by (simp add: Arg-zero angle-c-def)

lemma angle-c-neq0:⟨angle-c p q r ≠ 0 ⟹ p≠r ⟹
  unfolding angle-c-def
  by force

end
theory Complex-Triangles-Definitions

imports Complex-Angles

begin

```

3 Complex triangles

In this section we define triangles and derive some useful lemmas on congruent triangles, following the model [1]

```

locale congruent-ctriangle =
  fixes a1 b1 c1 :: complex and a2 b2 c2 :: complex
  assumes sides': cdist a1 b1 = cdist a2 b2 cdist a1 c1 = cdist a2 c2 cdist b1 c1
  = cdist b2 c2
  and angles': angle-c b1 a1 c1 = angle-c b2 a2 c2 ∨ angle-c b1 a1 c1 = −
    angle-c b2 a2 c2

```

```

angle-c a1 b1 c1 = angle-c a2 b2 c2 ∨ angle-c a1 b1 c1 = - angle-c a2 b2 c2
angle-c a1 c1 b1 = angle-c a2 c2 b2 ∨ angle-c a1 c1 b1 = - angle-c a2 c2 b2
begin

lemma sides:
cdist a1 b1 = cdist a2 b2 cdist a1 c1 = cdist a2 c2 cdist b1 c1 = cdist b2 c2
cdist b1 a1 = cdist a2 b2 cdist c1 a1 = cdist a2 c2 cdist c1 b1 = cdist b2 c2
cdist a1 b1 = cdist b2 a2 cdist a1 c1 = cdist c2 a2 cdist b1 c1 = cdist c2 b2
cdist b1 a1 = cdist b2 a2 cdist c1 a1 = cdist c2 a2 cdist c1 b1 = cdist c2 b2
using sides'
by (auto simp add: norm-minus-commute)

lemma angles:
angle-c b1 a1 c1 = angle-c b2 a2 c2 ∨ angle-c b1 a1 c1 = - angle-c b2 a2 c2
angle-c a1 b1 c1 = angle-c a2 b2 c2 ∨ angle-c a1 b1 c1 = - angle-c a2 b2 c2
angle-c a1 c1 b1 = angle-c a2 c2 b2 ∨ angle-c a1 c1 b1 = - angle-c a2 c2 b2
angle-c c1 a1 b1 = angle-c b2 a2 c2 ∨ angle-c c1 a1 b1 = - angle-c b2 a2 c2
angle-c c1 b1 a1 = angle-c a2 b2 c2 ∨ angle-c c1 b1 a1 = - angle-c a2 b2 c2
angle-c b1 c1 a1 = angle-c a2 c2 b2 ∨ angle-c b1 c1 a1 = - angle-c a2 c2 b2
angle-c b1 a1 c1 = angle-c c2 a2 b2 ∨ angle-c b1 a1 c1 = - angle-c c2 a2 b2
angle-c a1 b1 c1 = angle-c c2 b2 a2 ∨ angle-c a1 b1 c1 = - angle-c c2 b2 a2
angle-c a1 c1 b1 = angle-c b2 c2 a2 ∨ angle-c a1 c1 b1 = - angle-c b2 c2 a2
angle-c c1 a1 b1 = angle-c c2 a2 b2 ∨ angle-c c1 a1 b1 = - angle-c c2 a2 b2
angle-c c1 b1 a1 = angle-c c2 b2 a2 ∨ angle-c c1 b1 a1 = - angle-c c2 b2 a2
angle-c b1 c1 a1 = angle-c b2 c2 a2 ∨ angle-c b1 c1 a1 = - angle-c b2 c2 a2
using angles' angle-c-commute
by(simp-all add: angle-c-commute)(metis add.inverse-inverse angle-c-commute) +
end

```

```

end
theory Complex-Trigonometry

```

```

imports HOL-Analysis.Convex Complex-Angles Complex-Triangles-Definitions
begin

```

4 Complex trigonometry

In this section we add some trigonometric results, especially the law of sines

```

lemma ang-sin:
shows Im ((b-a)*cnj(c-a)) = cmod (c-a) *cmod (b-a) * sin (angle-c c a b)
proof -
have f0: sin (Arg (cnj c - cnj a)) = - sin (Arg (c - a))
by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus
sin-pi)

```

```

have f1: $\cos(\operatorname{Arg}(\operatorname{cnj} c - \operatorname{cnj} a)) = \cos(\operatorname{Arg}(c - a))$ 
  by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus)
then have  $\langle b - a = \operatorname{cmod}(b - a) * \operatorname{cis}(\operatorname{Arg}(b - a)) \wedge \operatorname{cnj}(c - a) = \operatorname{cmod}(c - a)$ 
  *  $\operatorname{cis}(\operatorname{Arg}(\operatorname{cnj}(c - a))) \rangle$ 
  using rcis-cmod-Arg rcis-def
  by (metis complex-mod-cnj)
then have  $\langle \operatorname{Im}(\operatorname{cis}(\operatorname{Arg}(b - a)) * \operatorname{cis}(\operatorname{Arg}(\operatorname{cnj}(c - a)))) = \sin(\operatorname{angle-c} c a b) \rangle$ 
  unfolding ang-vec-def angle-c-def using f1 f0 by(auto simp:cis.code sin-diff)
then have  $\langle \operatorname{Im}(\operatorname{cmod}(b - a) * \operatorname{cis}(\operatorname{Arg}(b - a)) * \operatorname{cmod}(c - a) * \operatorname{cis}(\operatorname{Arg}(\operatorname{cnj}(c - a)))) =$ 
   $\operatorname{cmod}(c - a) * \operatorname{cmod}(b - a) * \sin(\operatorname{angle-c} c a b) \rangle$ 
  by (metis Im-complex-of-real[of cmod(c - a)] Im-complex-of-real[of cmod(b - a)]
  [Im-mult-real[of cor(cmod(c - a))
  cis(Arg(cnj(c - a))) * (cor(cmod(b - a)) * cis(Arg(b - a)))]
  Im-mult-real[of cor(cmod(b - a)) cis(Arg(b - a)) * cis(Arg(cnj(c - a)))]]
  Re-complex-of-real[of cmod(c - a)] Re-complex-of-real[of cmod(b - a)]
  ab-semigroup-mult-class.mult-ac(1)[of cor(cmod(b - a)) cis(Arg(b - a))
  cis(Arg(cnj(c - a)))]
  ab-semigroup-mult-class.mult-ac(1)[of cis(Arg(cnj(c - a)))]
  cor(cmod(b - a)) * cis(Arg(b - a)) cor(cmod(c - a))]
  ab-semigroup-mult-class.mult-ac(1)[of cmod(c - a) cmod(b - a)
  Im(cis(Arg(b - a)) * cis(Arg(cnj(c - a))))]
  mult.commute[of cis(Arg(cnj(c - a))) cor(cmod(b - a)) * cis(Arg(b - a))]
  mult.commute[of cis(Arg(cnj(c - a))) * (cor(cmod(b - a)) * cis(Arg(b - a)))]
  cor(cmod(c - a))]
  mult.commute[of cor(cmod(b - a)) * cis(Arg(b - a)) * cor(cmod(c - a))
  cis(Arg(cnj(c - a)))]
  then show  $\langle \operatorname{Im}((b - a) * \operatorname{cnj}(c - a)) = \operatorname{cmod}(c - a) * \operatorname{cmod}(b - a) * \sin(\operatorname{angle-c} c a b) \rangle$ 
  by (metis <b - a = cor(cmod(b - a)) * cis(Arg(b - a))  $\wedge \operatorname{cnj}(c - a) =$ 
  cor(cmod(c - a)) * cis(Arg(cnj(c - a)))>
  ab-semigroup-mult-class.mult-ac(1))
qed

lemma ang-cos:
shows  $\operatorname{Re}((b - a) * \operatorname{cnj}(c - a)) = \operatorname{cmod}(c - a) * \operatorname{cmod}(b - a) * \cos(\operatorname{angle-c} c a b)$ 
proof -
have f0: $\langle \sin(\operatorname{Arg}(\operatorname{cnj} c - \operatorname{cnj} a)) = -\sin(\operatorname{Arg}(c - a)) \rangle$ 
by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus
  sin-pi)
have f1: $\langle \cos(\operatorname{Arg}(\operatorname{cnj} c - \operatorname{cnj} a)) = \cos(\operatorname{Arg}(c - a)) \rangle$ 
by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus)
then have f2: $\langle b - a = \operatorname{cmod}(b - a) * \operatorname{cis}(\operatorname{Arg}(b - a)) \wedge \operatorname{cnj}(c - a) = \operatorname{cmod}(c - a)$ 
  *  $\operatorname{cis}(\operatorname{Arg}(\operatorname{cnj}(c - a))) \rangle$ 

```

```

using rcis-cmod-Arg rcis-def
by (metis complex-mod-cnj)
then have ‹Re(cis (Arg (b-a)) * cis (Arg (cnj(c-a)))) = cos (angle-c c a b)›
  unfolding ang-vec-def angle-c-def using f1 f0 by(auto simp:cis.code cos-diff)
then have ‹Re (cmod (b-a) * cmod (c-a)* cis (Arg (b-a)) * cis (Arg (cnj
(c-a)))) =
  Re (cmod (c-a) *cmod (b-a)) * cos (angle-c c a b) ›
proof -
  have f1: cor (cmod (b - a) * cmod (c - a)) * cis (Arg (b - a)) * cis (Arg
(cnj (c - a)))
  = cor (cmod (b - a) * cmod (c - a)) * (cis (Arg (b - a)) * cis (Arg (cnj (c -
a))))
    using ab-semigroup-mult-class.mult-ac(1) by blast
  have cmod (b - a) * cmod (c - a) = cmod (c - a) * cmod (b - a)
    by argo
  then show ?thesis
    using f1 Im-complex-of-real Re-mult-real ‹Re (cis (Arg (b - a)) * cis (Arg
(cnj (c - a))) )
      = cos (angle-c c a b)› by presburger
  qed
  then show ‹Re ((b-a)*cnj(c-a)) = cmod (c-a) *cmod (b-a) * cos (angle-c c
a b)›
    by (metis Re-complex-of-real f2 mult.assoc mult.commute of-real-mult)

qed

```

lemma law-of-cosines':
assumes h: $\langle A \neq C \rangle \langle A \neq B \rangle$
shows $((cdist B C)^2 - (cdist A C)^2 - (cdist A B)^2) / (- 2 * (cdist A C) * (cdist A B)) = (\cos(\angle(C-A)(B-A)))$
using law-of-cosines[of B C A] h **by**(auto simp:field-simps)

lemma law-of-cosines'':
shows $(cdist A C)^2 = (cdist B C)^2 - (cdist A B)^2 + 2 * (cdist A C) * (cdist A B) * (\cos(\angle(C-A)(B-A)))$
using law-of-cosines[of B C A] **by**(auto simp:field-simps)

lemma law-of-cosines''':
shows $(cdist A B)^2 = (cdist B C)^2 - (cdist A C)^2 + 2 * (cdist A C) * (cdist A B) * (\cos(\angle(C-A)(B-A)))$
using law-of-cosines[of B C A] **by**(auto simp:field-simps)

4.1 The law of sines

theorem law-of-sines:
assumes h1: $\langle b \neq a \rangle \langle a \neq c \rangle \langle b \neq c \rangle$
shows $\sin(\angle-c a b c) * cdist b c = \sin(\angle-c c a b) * cdist a c$ (**is** ?A = ?B)
proof -

```

{fix a b c ::complex
  have f0:<sin (Arg (cnj c - cnj a)) = - sin (Arg (c - a))>
    by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus
        sin-pi)
  have f1:<cos (Arg (cnj c - cnj a)) = cos (Arg (c - a))>
    by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus)
  assume <b-a ≠ 0> <c-a ≠ 0>
  then have f2:<b-a = cmod (b-a) * cis (Arg (b-a)) ∧ cnj (c-a) = cmod
(c-a) * cis (Arg (cnj (c-a)))>
    using rcis-cmod-Arg rcis-def
    by (metis complex-mod-cnj)
  then have <Im(cis (Arg (b-a)) * cis (Arg (cnj(c-a)))) = sin (angle-c c a b)>
    unfolding ang-vec-def angle-c-def using f1 f0 by(auto simp:cis.code sin-diff)

  then have <Im (cmod (b-a) * cis (Arg (b-a)) * cmod (c-a) * cis (Arg (cnj
(c-a)))) =
      cmod (c-a) * cmod (b-a) * sin (angle-c c a b)>
    by (metis Im-complex-of-real[of cmod (c - a)] Im-complex-of-real[of cmod (b
- a)]]
      Im-mult-real[of cor (cmod (c - a))]
      cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg (b - a)))
      Im-mult-real[of cor (cmod (b - a)) cis (Arg (b - a)) * cis (Arg (cnj (c -
a)))]
      Re-complex-of-real[of cmod (c - a)] Re-complex-of-real[of cmod (b - a)]
      ab-semigroup-mult-class.mult-ac(1)[of cor (cmod (b - a)) cis (Arg (b -
a))]
      cis (Arg (cnj (c - a)))]
      ab-semigroup-mult-class.mult-ac(1)[of cis (Arg (cnj (c - a)))]
      cor (cmod (b - a)) * cis (Arg (b - a)) cor (cmod (c - a))]
      ab-semigroup-mult-class.mult-ac(1)[of cmod (c - a) cmod (b - a)
      Im (cis (Arg (b - a)) * cis (Arg (cnj (c - a))))]
      mult.commute[of cis (Arg (cnj (c - a))) cor (cmod (b - a)) * cis (Arg (b
- a))]
      mult.commute[of cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg
(b - a)))]
      cor (cmod (c - a))]
      mult.commute[of cor (cmod (b - a)) * cis (Arg (b - a)) * cor (cmod (c
- a))]
      cis (Arg (cnj (c - a)))]
    then have <Im ((b-a)*cnj(c-a)) = cmod (c-a) * cmod (b-a) * sin (angle-c
c a b)>
      using ang-sin by presburger
  }note ang=this
  have i2:<sin (angle-c c a b) = Im((b-a)*cnj(c-a))/ (cmod(b-a)*cmod(c-a))>
    using ang[of b a c] h1 by(auto)
  have <sin (angle-c a b c) = Im((c-b)*cnj(a-b)) / (cmod (c-b)*cmod (a-b))>
    using ang h1 by(auto)
  then have imp1:<cmod (c-b) * sin (angle-c a b c) = Im((c-b)*cnj(a-b)) / (cmod (a-b))>

```

```

    by auto
from i2 have imp2:⟨cmod (c-a) * sin (angle-c c a b) = Im((b-a)*cnj(c-a))/(
cmod(b-a)) ⟩
    by auto
show ?thesis using imp1 imp2 by(auto simp:field-simps norm-minus-commute)

qed

lemma law-of-sines': assumes h1:⟨b ≠ a⟩ ⟨a ≠ c⟩ ⟨b ≠ c⟩
shows sin (angle-c a b c) * cdist b a = sin (angle-c b c a) * cdist a c
proof -
{fix a b c ::complex
have f0:⟨sin (Arg (cnj c - cnj a)) = - sin (Arg (c - a))⟩
by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff neg-0-equal-iff-equal sin-minus
sin-pi)
have f1:⟨cos (Arg (cnj c - cnj a)) = cos (Arg (c - a))⟩
by (metis arg-cnj-not-pi arg-cnj-pi complex-cnj-diff cos-minus)
assume ⟨b-a ≠ 0⟩ ⟨c-a ≠ 0⟩
then have f2:⟨b-a = cmod (b-a) * cis (Arg (b-a)) ∧ cnj (c-a) = cmod
(c-a) * cis (Arg (cnj (c-a)))⟩
using rcis-cmod-Arg rcis-def
by (metis complex-mod-cnj)
then have ⟨Im(cis (Arg (b-a)) * cis (Arg (cnj(c-a)))) = sin (angle-c c a b))⟩
unfolding ang-vec-def angle-c-def using f0 f1 by(auto simp:cis.code sin-diff)

then have ⟨Im (cmod (b-a) * cis (Arg (b-a)) * cmod (c-a) * cis (Arg (cnj
(c-a)))) =
cmod (c-a) * cmod (b-a) * sin (angle-c c a b) ⟩
by (metis Im-complex-of-real[of cmod (c - a)] Im-complex-of-real[of cmod (b
- a)]
Im-mult-real[of cor (cmod (c - a))
cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg (b - a)))]
Im-mult-real[of cor (cmod (b - a)) cis (Arg (b - a)) * cis (Arg (cnj (c
- a)))]]
Re-complex-of-real[of cmod (c - a)] Re-complex-of-real[of cmod (b - a)]
ab-semigroup-mult-class.mult-ac(1)[of cor (cmod (b - a)) cis (Arg (b -
a))
cis (Arg (cnj (c - a)))]
ab-semigroup-mult-class.mult-ac(1)[of cis (Arg (cnj (c - a)))]
cor (cmod (b - a)) * cis (Arg (b - a)) cor (cmod (c - a))
ab-semigroup-mult-class.mult-ac(1)[of cmod (c - a) cmod (b - a)
Im (cis (Arg (b - a)) * cis (Arg (cnj (c - a))))]
mult.commute[of cis (Arg (cnj (c - a))) cor (cmod (b - a)) * cis (Arg (b
- a))]
mult.commute[of cis (Arg (cnj (c - a))) * (cor (cmod (b - a)) * cis (Arg
(b - a)))]
cor (cmod (c - a))
mult.commute[of cor (cmod (b - a)) * cis (Arg (b - a)) * cor (cmod (c
- a))]
```

```

    cis (Arg (cnj (c - a)))]
then have <Im ((b-a)*cnj(c-a)) = cmod (c-a) *cmod (b-a) * sin (angle-c
c a b)>
    using ang-sin by presburger
}note ang=this
have i2:<sin (angle-c b c a) = Im((a-c)*cnj(b - c))/ (cmod(b-c)*cmod(a-c))>
    using ang[of a c b] h1 by(auto)
have <sin (angle-c a b c) = Im((c-b)*cnj(a-b)) / (cmod (c-b)*cmod (a-b))>
    using ang h1 by(auto)
then have imp1:<cmod (a-b) * sin (angle-c a b c) = Im((c-b)*cnj(a-b)) /
(cmod (c-b)) >
    by auto
from i2 have imp2:<cmod (a-c) * sin (angle-c b c a) = Im((a-c)*cnj(b-c))/(
cmod(b-c)) >
    by auto
show ?thesis using imp1 imp2
    by (metis cdist-commute h1(1,2,3) law-of-sines)
qed

```

```

lemma ang-pos-pos:<q≠p ⇒ p≠ r ⇒ r ≠ q ⇒ angle-c q r p ≥ 0 ⇒ angle-c
r p q ≥ 0>
    using law-of-sines'[of r q p]
    by (smt (verit) ang-vec-bounded angle-c-def cdist-def mult-neg-pos mult-nonneg-nonneg
right-minus-eq sin-ge-zero sin-gt-zero sin-minus zero-less-norm-iff)

```

```

lemma cmod-pos:<cmod a ≥ 0>
    by simp

```

```

lemma ang-neg-neg:<q≠p ⇒ p≠ r ⇒ r ≠ q ⇒ angle-c q r p < 0 ⇒ angle-c
r p q < 0>
proof -
    assume <q≠p> <p≠ r> <r ≠ q> <angle-c q r p < 0>
    then have <sin (angle-c q r p) < 0>
        using ang-c-in
        by (metis ang-vec-def angle-c-def canon-ang(1) minus-less-iff neg-0-less-iff-less
sin-gt-zero sin-minus)
    then have <sin (angle-c q r p) * cdist r q < 0>
        using <r ≠ q> mult-neg-pos by fastforce
    from law-of-sines'[of r q p] have <sin (angle-c r p q)<0>
        by (metis <p ≠ r> <q ≠ p> <r ≠ q> <sin (angle-c q r p) * cdist r q < 0>
cdist-def linorder-not-less mult-less-0-iff norm-ge-zero)
    then show ?thesis
        using ang-c-in[of r p q]
        by (metis ang-vec-def angle-c-def canon-ang(2) linorder-not-less sin-ge-zero)
qed

```

```

lemma collinear-sin-neq-0:

```

```

< $\neg \text{collinear } a2 b2 c2 \implies \sin(\text{angle-}c\ a2\ c2\ b2) \neq 0$ >
unfolding collinear-def angle-c-def
by (metis Im-complex-div-eq-0 ang-sin angle-c-def collinear-def
collinear-sym1 collinear-sym2' mult-eq-0-iff)

```

lemma *collinear-sin-neq-pi*:

```

< $\neg \text{collinear } a2 b2 c2 \implies \sin(\text{angle-}c\ a2\ c2\ b2) \neq \pi$ >
unfolding collinear-def angle-c-def
by (metis add-cancel-right-left dual-order.antisym dual-order.trans le-add-same-cancel1
linordered-nonzero-semiring-class.zero-le-one one-add-one one-neq-zero pi-ge-two sin-le-one)

```

lemma *collinear-iff*:

```

assumes  $a \neq b \wedge b \neq c \wedge c \neq a$ 
shows  $\langle \text{collinear } a\ b\ c \longleftrightarrow (\text{angle-}c\ a\ b\ c = \pi \vee \text{angle-}c\ a\ b\ c = 0) \rangle$ 
apply(rule iffI)
using assms unfolding collinear-def using collinear-angle collinear-def apply(fastforce)
by (metis collinear-def collinear-sin-neq-0 collinear-sym1 sin-pi sin-zero)

```

definition $\langle \text{innerprod } a\ b \equiv \text{cnj } a * b \rangle$

lemma *left-lin-innerprod*: $\langle \text{innerprod } (x + y)\ z = \text{innerprod } x\ z + \text{innerprod } y\ z \rangle$

```

unfolding innerprod-def
by (simp add: mult.commute ring-class.ring-distrib(1))

```

lemma *right-lin-innerprod*: $\langle \text{innerprod } x\ (y+z) = \text{innerprod } x\ y + \text{innerprod } x\ z \rangle$

```

unfolding innerprod-def
by (simp add: ring-class.ring-distrib(1))

```

lemma *leftlin-innerprod*: $\langle \text{innerprod } x\ (t*y) = t * \text{innerprod } x\ y \rangle$

```

unfolding innerprod-def by(auto)

```

lemma *rightsesqlin-innerprod*: $\langle \text{innerprod } (t*x)\ (y) = \text{cnj } t * \text{innerprod } x\ y \rangle$

```

unfolding innerprod-def by(auto)

```

lemma *norm-eq-csqrt-inner*: $\langle \text{norm } x = \text{csqrt } (\text{innerprod } x\ x) \rangle$

```

using complex-mod-sqrt-Re-mult-cnj innerprod-def by force

```

lemma *abs2-eq-inner*: $\langle \text{abs } (\text{innerprod } x\ y)^2 = \text{innerprod } x\ y * \text{cnj } (\text{innerprod } x\ y) \rangle$

```

unfolding innerprod-def abs-complex-def apply(rule complex-eqI)
by (metis comp-apply complex-mult-cnj-cmod of-real-power)

```

```

lemma complex-add-inner-cnj:⟨t*innerprod x y + cnj (t*innerprod x y) = 2* Re
(t*innerprod x y)⟩
  using complex-add-cnj by blast

lemma Re-innerprod-inner:⟨Re (innerprod (a-b) (c-b)) = (a-b)•(c-b)⟩
  unfolding innerprod-def inner-complex-def by(auto simp:field-simps)

lemma angle-c-arccos-pos:
  assumes h:⟨a≠b ∧ b≠c ∧ angle-c a b c ≥ 0⟩
  shows ⟨angle-c a b c = arccos ((Re (innerprod (a-b) (c - b)))/(cmod(a-b)*cmod(c-b)))⟩
proof -
  have ⟨Re ((a-b)*cnj(c-b)) = (Re (innerprod (a-b) (c - b)))⟩
    unfolding innerprod-def by(auto simp:field-simps)
  then have ⟨((Re (innerprod (a-b) (c - b)))/(cmod(a-b)*cmod(c-b))) = cos
(angle-c a b c))⟩
    by (metis (no-types, lifting) ang-cos angle-c-commute cos-minus divisors-zero
h(1) mult.commute nonzero-mult-div-cancel-left norm-eq-zero right-minus-eq)
  then show ?thesis
    using ang-vec-bounded angle-c-def arccos-cos h(2) by presburger
qed

lemma angle-c-arccos-neg:
  assumes h:⟨a≠b ∧ b≠c ∧ angle-c a b c ≤ 0 ⟩
  shows ⟨- angle-c a b c = arccos ((Re (innerprod (a-b) (c - b)))/(cmod(a-b)*cmod(c-b)))⟩
proof -
  have ⟨Re ((a-b)*cnj(c-b)) = (Re (innerprod (a-b) (c - b)))⟩
    unfolding innerprod-def by(auto simp:field-simps)
  then have ⟨((Re (innerprod (a-b) (c - b)))/(cmod(a-b)*cmod(c-b))) = cos
(angle-c a b c))⟩
    by (metis (no-types, lifting) ang-cos angle-c-commute cos-minus divisors-zero
h(1) mult.commute nonzero-mult-div-cancel-left norm-eq-zero right-minus-eq)
  then show ?thesis
    using ang-vec-bounded angle-c-def arccos-cos h(2)
    by (metis arccos-cos2 less-le-not-le)
qed

end
theory Third-Unity-Root

imports Complex-Angles

begin

```

5 Third unity root

In this section we prove some basic properties of the third unity root j.

```
lemma root-unity-3: ⟨(z::complex)^3 - 1 = 0 ⟺ (z = cis (2*pi/3) ∨ z=1 ∨ z
```

```

= cis (4*pi/3))>
proof(rule iffI)
  have < cis (2 * pi * real xa / 3) ≠ 1 =>
    cis (2 * pi * real xa / 3) ≠ cis (2 * pi / 3)
  ⟹ xa < 3 => cis (2 * pi * real xa / 3) = cis (4 * pi / 3) ›
    for xa
  proof(rule ccontr)
    assume h:< cis (2 * pi * real xa / 3) ≠ 1 ›
    < cis (2 * pi * real xa / 3) ≠ cis (2 * pi / 3)›
    < xa < 3› < cis (2 * pi * real xa / 3) ≠ cis (4 * pi / 3)›
    then have <xa = 1 ∨ xa = 0›
      using less-Suc-eq numeral-3-eq-3 by fastforce
    then show False
      using h(1) h(2) by auto
    qed
    then have f0:<(λk. cis (2 * pi * real k / real 3)) ` {..<3} = {1, cis(2*pi/3), cis (4*pi/3)}›
      unfolding image-def by(auto simp:image-def intro:bexI[where x=0]
        bexI[where x=1] bexI[where x=2])
    have <{z. z^3 = 1} = {1, cis(2*pi/3), cis (4*pi/3)}›
      apply(insert bij-betw-roots-unity[of 3, simplified])
      apply(frule bij-betw-imp-surj-on)
      using f0 by auto
    then show <z ^ 3 - 1 = 0 => z = cis (2 * pi / 3) ∨ z = 1 ∨ z = cis (4 * pi
    / 3)›
      unfolding bij-betw-def inj-on-def by(auto)
      show <z = cis (2 * pi / 3) ∨ z = 1 ∨ z = cis (4 * pi / 3) => z ^ 3 - 1 = 0
    ›
      using cis-2pi cis-multiple-2pi[of 2] by(auto simp:DeMoivre field-simps)
    qed

definition root3 where <root3≡cis(2*pi/3)›

lemma root3-eq-0:<1 + root3 + root3^2 = 0›
proof –
  have <(z::complex)^3 - 1 = (z - 1)*(1 + z + z^2)› for z
    by(auto simp:field-simps power2-eq-square power3-eq-cube)
  moreover have <root3^3 - 1 = 0›
    using root-unity-3 unfolding root3-def by blast
  moreover have <root3 - 1 ≠ 0›
    by(simp add:cis.code Complex-eq-1 cos-120 root3-def)
  ultimately show ?thesis
    by simp
qed

lemma root-unity-carac:
  assumes h1:<a1*a2*a3 = j› <1-a1*a2 ≠ 0 ∧ 1-a2*a3≠0 ∧ 1-a1*a3 ≠ 0›
  <j=root3›
  and R : <R=(a1 * b2 + b1) / (1 - a1 * a2)› (is <R=?R›)

```

```

and P :<P=(a2*b3 + b2)/(1-a2*a3)> (is <P=?P>)
and Q:<Q=(a3*b1 + b3)/(1-a3*a1)> (is <Q=?Q>)
shows <(a1^2+a1+1)*b1 + a1^3*(a2^2+a2+1)*b2 + a1^3*a2^3*(a3^2+a3+1)*b3
=
- j*a1^2*a2*(a1-j)*(a2-j)*(?R +j*?P +j^2*?Q)>
proof -
have simp-rules-for-eq:<1-a1*a2 ≠ 0 ∧ 1-a2*a3≠0 ∧ 1-a1*a3 ≠ 0> <a1*a2*a3
= j>
<1 + j +j^2 = 0> <j*j*j = 1> <(1-a1*a2)*a3 = (a3-j)> <(1-a2*a3)*a1 =
(a1-j)> <(1-a1*a3)*a2 = (a2-j)>
using h1 root3-eq-0 root-unity-3 unfolding root3-def
by(auto simp:power3-eq-cube mult.left-commute mult.commute right-diff-distrib)

have y:<(- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * ((1 - a1 * a3) * a2) * ((1
- a1 * a2) * a3) *
(a1 * a2) + - j * a1^2 * a2 * ((1 - a2 * a3) * a1) * ((1 - a1 * a3) * a2) *
((1 - a1 * a2) * a3) * b1) /
(1 - a1 * a2) = (- j * a1^2 * a2 * ((1 - a3 * a2) * a1) * ((1 - a1 * a3) *
a2) * ( a3) * (a1 * b2)
+ - j * a1^2 * a2 * ((1 - a3 * a2) * a1) * ((1 - a1 * a3) * a2) * (a3) *
b1)>
apply(subst add-divide-distrib)
using simp-rules-for-eq(1)
by(simp add: mult.commute)
have y2:<(- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * ((1 - a1 * a3) * a2) * ((1
- a1 * a2) * a3) * (j * (a2 * b3)) +
- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * ((1 - a1 * a3) * a2) * ((1 - a1 *
a2) * a3) * (j * b2)) /
(1 - a2 * a3) = (- j * a1^2 * a2 * (a1) * ((1 - a1 * a3) * a2) * ((1 - a1
* a2) * a3) * (j * (a2 * b3)) +
- j * a1^2 * a2 * (a1) * ((1 - a1 * a3) * a2) * ((1 - a1 * a2) * a3) * (j *
b2))>
apply(subst add-divide-distrib)
using simp-rules-for-eq(1) by auto
have y3:<(- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * ((1 - a1 * a3) * a2) * ((1
- a1 * a2) * a3) * (j^2 * (a3 * b1)) +
- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * ((1 - a1 * a3) * a2) * ((1 - a1 *
a2) * a3) * (j^2 * b3)) /
(1 - a3 * a1) =
(- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * (a2) * ((1 - a1 * a2) * a3) * (j^2 *
(a3 * b1)) +
- j * a1^2 * a2 * ((1 - a2 * a3) * a1) * (a2) * ((1 - a1 * a2) * a3) * (j^2 *
b3))>
apply(subst add-divide-distrib)
by(smt (verit, best) mult.commute nonzero-mult-div-cancel-left times-divide-eq-left
simp-rules-for-eq(1))
have ko:<a*(b*(c)) = a*b*c> for a b c::complex
by auto
have ko':<a +(b +(c)) = a+b+c> for a b c::complex

```

```

by auto
have *:( $a_1 * a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_2 * a_3 * a_3 * a_3 * b_1$ 
=  $a_1 * a_1 * a_2 * a_3 * b_1$ )
  using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
have **:( $a_1 * a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_2 * a_3 * a_3 * a_3 * b_3$ 
=  $a_1 * a_1 * a_2 * a_2 * b_3$ )
  using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
have ***:( $a_1 * a_1 * a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_2 * a_3 * a_3 * a_3 * a_3 * b_3$ 
=  $a_1 * a_1 * a_2 * a_3 * b_3$ )
  using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
have ****:( $a_1 * a_1 * a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_2 * a_3 * a_3 * a_3 * a_3 * b_3$ 
=  $a_1 * a_1 * a_1 * a_2 * a_3 * b_3$ )
  using simp-rules-for-eq by (auto simp add: mult.assoc mult.left-commute)
have fin:( $a_1 * b_1 + a_1 * a_1 * a_1 * a_2 * b_2 + a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_3 * b_2 +$ 
 $a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_3 * b_3 + a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_3 * a_3 * b_1 +$ 
 $a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_3 * a_3 * b_2 + a_1 * a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_2 * a_3 * a_3 * b_1 +$ 
 $a_1 * a_1 * a_1 * a_1 * a_2 * a_2 * a_2 * a_3 * a_3 * b_3 + a_1 * a_1 * a_2 * a_2 * a_3 * b_1 = a_1 * b_1$ 
+  $j * a_1 * b_1 + j^2 * a_1 * b_1 +$ 
 $a_1 * a_1 * a_1 * a_2 * b_2 + a_1 * a_1 * a_1 * a_2 * b_2 * j + a_1 * a_1 * a_1 * a_2 * b_2 * j^2 + a_1 * a_1 * a_2 * a_2 * b_3$ 
+  $a_1 * a_1 * a_2 * a_2 * b_3 * j +$ 
 $a_1 * a_1 * a_2 * a_2 * b_3 * j^2$ )
  using simp-rules-for-eq by (auto simp: algebra-simps power2-eq-square power3-eq-cube)
then have fin':... = ( $a_1 * b_1 + a_1 * a_1 * a_1 * a_2 * b_2 + a_1 * a_1 * a_2 * a_2 * b_3$ ) * (1 +  $j + j^2$ )
  by (auto simp: algebra-simps power2-eq-square power3-eq-cube)
then show graal:( $(a_1^2 + a_1 + 1) * b_1 + a_1^3 * (a_2^2 + a_2 + 1) * b_2 + a_1^3 * a_2^3 * (a_3^2 + a_3 + 1) * b_3$ 
=  $-j * a_1^2 * a_2 * (a_1 - j) * (a_2 - j) * (a_3 - j) * (?R + j * ?P + j^2 * ?Q)$ )
  apply(simp only: simp-rules-for-eq(5-7)[symmetric])
  apply(simp only:y y2 y3 distrib-left distrib-right times-divide-eq-left times-divide-eq-right)

apply(simp only:simp-rules-for-eq algebra-simps del:mult.assoc mult.commute)

apply(simp only:ko ko')
using simp-rules-for-eq apply(auto simp add: power2-eq-square power3-eq-cube
algebra-simps)
  apply(simp only:ko ko') apply(subst *) apply(subst **) apply(subst ***)
apply(simp)
  apply(simp only:ko ko') apply(subst fin) apply(subst fin')
  using simp-rules-for-eq by fastforce
qed

end
theory Complex-Triangles

imports Complex-Trigonometry Third-Unity-Root

begin

```

lemma *similar-triangles'*:

assumes $h:a \neq 0 \wedge b \neq 0 \wedge c \neq 0 \wedge a' \neq 0 \wedge b' \neq 0 \wedge c' \neq 0$
and $h1:\langle\angle(-a) b = \angle(-a') b'\rangle \langle\angle(-b') c' = \angle(-b) c\rangle$
shows $\langle\angle(-c) a = \angle(-c') a'\rangle$

proof –

have $\langle|\angle(-a) b + \angle(-b) c + \angle(-c) a| = pi\rangle$
using *angle-sum-triangle'* h **by** *auto*
also have $\langle|\angle(-a') b' + \angle(-b') c' + \angle(-c') a'| = pi\rangle$
using *angle-sum-triangle'* h **by** *auto*
also have $\langle|\angle(-a') b' + \angle(-b') c' + \angle(-c) a| = pi\rangle$
using $\langle|\angle(-a) b + \angle(-b) c + \angle(-c) a| = pi\rangle$ $h1(1,2)$ **by** *auto*
then have $\langle|\angle(-c) a| = |\angle(-c') a'|\rangle$
by (*smt (verit)* $\langle|\angle(-a') b' + \angle(-b') c' + \angle(-c') a'| = pi\rangle$ *add.inverse-inverse ang-vec-bounded ang-vec-opposite1 ang-vec-opposite2 ang-vec-opposite-opposite ang-vec-plus-pi1 ang-vec-plus-pi2 canon-ang-diff canon-ang-id h neg-0-equal-iff-equal*)
ultimately show *?thesis*
by (*metis ang-vec-bounded canon-ang-id*)
qed

lemma *similar-triangles*:

assumes $h:a \neq b \wedge b \neq c \wedge a \neq c \wedge a' \neq b' \wedge b' \neq c' \wedge c' \neq a'$
and $h1:\langle\angle(c-a) (b-a) = \angle(c'-a') (b'-a')\rangle \langle\angle(a-b) (c-b) = \angle(a'-b') (c'-b')\rangle$
shows $\langle\angle(b-c) (a-c) = \angle(b'-c') (a'-c')\rangle$

proof –

have $\langle a-b \neq 0 \rangle \langle b-c \neq 0 \rangle \langle a-c \neq 0 \rangle \langle a'-b' \neq 0 \rangle \langle b'-c' \neq 0 \rangle \langle c'-a' \neq 0 \rangle$
using h **by** *auto*
then show *?thesis*
by (*smt (z3) diff-numeral-special(12)* $h1(1,2)$ *minus-diff-eq similar-triangles'*)
qed

lemma *similar-triangles-c*:

assumes $h:a \neq b \wedge b \neq c \wedge a \neq c \wedge a' \neq b' \wedge b' \neq c' \wedge c' \neq a'$
and $h1:\langle\text{angle-c } a b = \text{angle-c } c' a' b'\rangle \langle\text{angle-c } a b c = \text{angle-c } a' b' c'\rangle$
shows $\langle\text{angle-c } b c a = \text{angle-c } b' c' a'\rangle$

proof –

have $\langle a-b \neq 0 \rangle \langle b-c \neq 0 \rangle \langle a-c \neq 0 \rangle \langle a'-b' \neq 0 \rangle \langle b'-c' \neq 0 \rangle \langle c'-a' \neq 0 \rangle$
using h **by** *auto*
then show *?thesis*
unfolding *angle-c-def*
by (*metis angle-c-def h h1(1,2) similar-triangles*)
qed

lemmas *congruent-ctriangleD = congruent-ctriangle.sides congruent-ctriangle.angles*

lemma congruent-ctrianglessss:

assumes $h:a \neq b \wedge b \neq c \wedge a \neq c$

and $h1:\langle cmod(b-a) = cmod(b'-a') \rangle \wedge \langle cmod(b-c) = cmod(b'-c') \rangle \wedge \langle cmod(c-a) = cmod(c'-a') \rangle$

shows $\langle \text{congruent-ctriangle } a\ b\ c\ a'\ b'\ c' \rangle$

proof –

{fix $c\ b\ a\ c'\ b'\ a'$

assume $h:a \neq b \wedge b \neq c \wedge a \neq c$

and $h1:\langle cmod(b-a) = cmod(b'-a') \rangle \wedge \langle cmod(b-c) = cmod(b'-c') \rangle \wedge \langle cmod(c-a) = cmod(c'-a') \rangle$

then have $h':\langle a' \neq b' \wedge b' \neq c' \wedge c' \neq a' \rangle$

by auto

have $\langle \cos(\angle(c-a)(b-a)) = ((cdist b\ c)^2 - (cdist a\ c)^2 - (cdist a\ b)^2) / (-2 * (cdist a\ c) * (cdist a\ b)) \rangle$

using law-of-cosines' h by auto

moreover have $\langle \cos(\angle(c'-a')(b'-a')) = ((cdist b'\ c')^2 - (cdist a'\ c')^2 - (cdist a'\ b')^2) / (-2 * (cdist a'\ c') * (cdist a'\ b')) \rangle$

using law-of-cosines' h h' by auto

ultimately have $f0:\langle \cos(\angle(c-a)(b-a)) = \cos(\angle(c'-a')(b'-a')) \rangle$

by (metis cdist-def h1(1,2,3) norm-minus-commute)

have $\langle \angle(c-a)(b-a) \in \{-pi..pi\} \rangle \wedge \langle \angle(c'-a')(b'-a') \in \{-pi..pi\} \rangle$

unfolding ang-vec-def

using canon-ang(1,2) less-eq-real-def by auto

with $f0$ have $\langle \angle(c-a)(b-a) = \angle(c'-a')(b'-a') \vee \angle(c-a)(b-a) = -\angle(c'-a')(b'-a') \rangle$

unfolding ang-vec-def

by (smt (verit) arccos-cos2 arccos-minus-1 arccos-unique canon-ang(1,2))}note dem=this

then show ?thesis

by (metis (mono-tags, lifting) angle-c-def cdist-def congruent-ctriangle-def h h1(1) h1(2) h1(3) norm-minus-commute)

qed

lemma congruent-ctriangleIsss:

assumes $h:a \neq b \wedge b \neq c \wedge a \neq c$

and $h1:\langle cdist a\ b = cdist a'\ b' \rangle \wedge \langle cdist b\ c = cdist b'\ c' \rangle \wedge \langle cdist a\ c = cdist a'\ c' \rangle$

shows $\langle \text{congruent-ctriangle } a\ b\ c\ a'\ b'\ c' \rangle$

using cdist-def congruent-ctrianglessss h1(1,2,3) h

by (metis dist-commute dist-complex-def)

lemmas congruent-ctrianglesss = congruent-ctriangled[*OF* congruent-ctriangleIsss]

lemma isosceles-triangles:

assumes $\langle cdist a\ b = cdist b\ c \rangle$

shows $\langle \text{angle-c } b\ c\ a = \text{angle-c } b\ a\ c \vee \text{angle-c } b\ c\ a = -\text{angle-c } b\ a\ c \rangle$

by (metis assms cdist-commute cdist-def congruent-ctriangle-sss(14) norm-eq-zero right-minus-eq)

lemma *non-collinear-independant*: $\neg \text{collinear } a b c \implies a \neq b \wedge b \neq c \wedge a \neq c$
using *collinear-ex-real* **by** *force*

lemma *congruent-ctriangleI-sas*:
assumes $a1 \neq b1 \wedge b1 \neq c1 \wedge a1 \neq c1$
assumes $h1:\text{cdist } a1 b1 = \text{cdist } a2 b2$
assumes $h2:\text{cdist } b1 c1 = \text{cdist } b2 c2$
assumes $h3:\text{angle-c } a1 b1 c1 = \text{angle-c } a2 b2 c2 \vee \text{angle-c } a1 b1 c1 = -\text{angle-c } a2 b2 c2$
shows *congruent-ctriangle* $a1 b1 c1 a2 b2 c2$
proof (*rule congruent-ctriangleI-sss*)
have $h1':\text{cdist } a1 b1 = \text{cdist } b2 a2$
using *cdist-commute* $h1$ **by** *auto*
show $\text{cdist } a1 c1 = \text{cdist } a2 c2$
proof (*rule power2-eq-imp-eq*)
show $(\text{cdist } a1 c1)^2 = (\text{cdist } a2 c2)^2$
apply (*insert law-of-cosines* [of $a1 c1 b1$] *law-of-cosines* [of $a2 c2 b2$])
apply (*subst* (*asm*) (1 2) $h2$) **apply** (*subst* (*asm*) (1 2) $h1$ [*symmetric*])
using *assms* **unfolding** *angle-c-def*
by (*smt (verit)* *congruent-ctriangle-sss* (7) *angle-c-commute* *angle-c-def cos-minus*)
qed simp-all
show $a1 \neq b1 \wedge b1 \neq c1 \wedge a1 \neq c1$
 $\text{cdist } a1 b1 = \text{cdist } a2 b2$
 $\text{cdist } b1 c1 = \text{cdist } b2 c2$ **using** *assms* **cdist-commute** **by** *auto*
qed

lemmas *congruent-ctriangle-sas* = *congruent-ctriangleD* [*OF congruent-ctriangleI-sas*]

lemma *congruent-ctriangleI-aas*:
assumes $h1:\text{angle-c } a1 b1 c1 = \text{angle-c } a2 b2 c2$
assumes $h2:\text{angle-c } b1 c1 a1 = \text{angle-c } b2 c2 a2$
assumes $h3:\text{cdist } a1 b1 = \text{cdist } a2 b2$
assumes $h4:\neg \text{collinear } a1 b1 c1 \neg \text{collinear } a2 b2 c2$
shows *congruent-ctriangle* $a1 b1 c1 a2 b2 c2$
proof (*rule congruent-ctriangleI-sas*)
from $h4$ **have** $\text{neq}: a1 \neq b1$ **unfolding** *collinear-def* **by** *auto*
with *assms*(3) **have** $\text{neq}': a2 \neq b2$ **by** *auto*
have $h0:(a1 \neq b1) \wedge (b1 \neq c1) \wedge (a1 \neq c1)$
using *assms*(4) *non-collinear-independant* **by** *auto*
have $h0':(a2 \neq b2) \wedge (b2 \neq c2) \wedge (a2 \neq c2)$
using *assms*(5) *non-collinear-independant* **by** *auto*
have $A:\text{angle-c } c1 a1 b1 = \text{angle-c } c2 a2 b2$ **using** *neq neq'* *assms*
apply (*insert angle-sum-triangle-c* [of $a1 b1 c1$] *angle-sum-triangle-c* [of $a2 b2 c2$])

```

c2])
  by (metis h0 assms(5) h1 h2 non-collinear-independant similar-triangles-c)
  have <-∠(b1 - c1) (a1 - c1) = -∠(b2 - c2) (a2 - c2)>
    using h2 unfolding angle-c-def by auto
  then have A1:<angle-c a1 c1 b1 = angle-c a2 c2 b2>
    using h2 unfolding angle-c-def
    apply(cases <angle-c a1 c1 b1 = pi>)
      apply (metis ang-vec-def angle-c-def canon-ang-uminus-pi minus-diff-eq)
      by (metis ang-vec-sym ang-vec-sym-pi)
    have <∠(b1 - a1) (c1 - a1) = ∠(b2 - a2) (c2 - a2)>
      by (metis ang-vec-sym ang-vec-sym-pi angle-c-def A)
    have sine1:<sin (angle-c a1 c1 b1) * cdist c1 b1 = sin (angle-c b1 a1 c1) * cdist
a1 b1>
      by (metis h0 law-of-sines)
    have sine2:<sin (angle-c a2 c2 b2) * cdist c2 b2 = sin (angle-c b2 a2 c2) * cdist
a2 b2>
      by (metis assms(5) law-of-sines non-collinear-independant)
    have <¬collinear a2 b2 c2 ⇒ sin (angle-c a2 c2 b2) ≠ 0>
      unfolding collinear-def angle-c-def
      using angle-c-def collinear-sin-neq-0 h4(2) by presburger
    have h3':<cdist b1 a1 = cdist b2 a2>
      using h3 cdist-commute by auto
    have A2:<angle-c b1 a1 c1 = angle-c b2 a2 c2>
      using <∠(b1 - a1) (c1 - a1) = ∠(b2 - a2) (c2 - a2)> angle-c-def by auto
    have < cdist c1 b1 = sin (angle-c b2 a2 c2) * cdist a2 b2 / sin (angle-c a2 c2
b2) ∧
      cdist c2 b2 = sin (angle-c b2 a2 c2) * cdist a2 b2 / sin (angle-c a2 c2 b2) >
      apply(cases <sin (angle-c a2 c2 b2) = 0>)
        apply (simp add:
          ¬ Elementary-Complex-Geometry.collinear a2 b2 c2 ⇒ sin (angle-c a2 c2
b2) ≠ 0)
          ¬ Elementary-Complex-Geometry.collinear a2 b2 c2)
        by (metis A1 A2 h3 nonzero-mult-div-cancel-left sine1 sine2)
    then show cdist b1 c1 = cdist b2 c2
      using cdist-commute by auto
    show a1 ≠ b1 ∧ b1 ≠ c1 ∧ a1 ≠ c1
      cdist a1 b1 = cdist a2 b2
      angle-c a1 b1 c1 = angle-c a2 b2 c2 ∨ angle-c a1 b1 c1 = - angle-c a2 b2 c2
      using assms h0 neq by auto
qed

```

lemmas congruent-ctrangle-aas = congruent-ctrangleD[*OF* congruent-ctrangleI-aas]

lemma congruent-ctrangleI-asa:

```

assumes angle-c a1 b1 c1 = angle-c a2 b2 c2
assumes cdist a1 b1 = cdist a2 b2
assumes h0:angle-c b1 a1 c1 = angle-c b2 a2 c2
assumes h4:¬collinear a1 b1 c1 ¬collinear a2 b2 c2

```

```

shows congruent-ctriangle a1 b1 c1 a2 b2 c2
proof (rule congruent-ctriangleI-aas)
  from assms have neq: a1 ≠ b1 a2 ≠ b2 unfolding collinear-def by auto
  show angle-c b1 c1 a1 = angle-c b2 c2 a2
    apply (rule similar-triangles-c)
    using assms(4,5) non-collinear-independant apply auto[1]
    using h0 unfolding angle-c-def
    apply (metis ang-vec-sym ang-vec-sym-pi)
    using angle-c-def assms(1) by auto
qed fact+

```

lemmas congruent-ctriangle-asa = congruent-ctriangleD[*OF* congruent-ctriangleI-asa]

```

lemma orientation-respect-letter-order:
  assumes angle-c a b c = angle-c b a c ∨ collinear a b c
  shows False
  by (smt (verit, ccfv-threshold) angle-c-commute assms(1)
       assms(2) collinear-sin-neq-0 collinear-sym1 similar-triangles-c sin-pi sin-zero)

```

```

lemma isoscele-iff-pr-cis-qr:
  assumes h': q ≠ r
  shows ⟨(cdist q r = cdist p r) ↔ (p - r) = cis(angle-c q r p) * (q - r)⟩
proof (rule iffI)
  assume h: ⟨(cdist q r = cdist p r)⟩
  then have f0: ⟨cmod(p - r) = cmod(q - r)⟩
    by (simp add: norm-minus-commute)
  have ⟨(p - r) / (q - r) = (p - r) * cnj(q - r) / cmod(q - r) ^ 2⟩
    using complex-div-cnj by blast
  have f1: ⟨Re((p - r) * cnj(q - r)) = cmod(p - r) * cmod(q - r) * cos(angle-c q r p)⟩
    by (metis (no-types, lifting) ang-cos cdist-commute cdist-def h(1))
  have f2: ⟨Im((p - r) * cnj(q - r)) = cmod(p - r) * cmod(q - r) * sin(angle-c q r p)⟩
    using ang-sin h(1) by auto
  then have ⟨(p - r) * cnj(q - r) / cmod(q - r) ^ 2 = cmod(p - r) * cmod(q - r) * cis(angle-c q r p) / cmod(q - r) ^ 2⟩
    apply(intro complex-eqI)
    using f1 f2 by auto
  then have f3: ⟨(p - r) / (q - r) = cmod(p - r) * cmod(q - r) * cis(angle-c q r p) / cmod(q - r) ^ 2⟩
    by (simp add: ⟨(p - r) / (q - r) = (p - r) * cnj(q - r) / cor((cmod(q - r) ^ 2))⟩)
  have ⟨(p - r) / (q - r) = cis(angle-c q r p)⟩
    apply(subst f3) apply(subst f0) by(simp add:h h' power2-eq-square)
  then show ⟨p - r = cis(angle-c q r p) * (q - r)⟩
    using divide-eq-eq h h' by force
next
assume ⟨p - r = cis(angle-c q r p) * (q - r)⟩
have ⟨(p - r) / (q - r) = (p - r) * cnj(q - r) / cmod(q - r) ^ 2⟩

```

```

using complex-div-cnj by blast
have f1:⟨Re ((p-r)*cnj(q-r)) = cmod(p-r)*cmod(q-r)*cos(angle-c q r p)⟩
  using ang-cos by fastforce
have f2:⟨Im ((p-r)*cnj(q-r)) = cmod(p-r)*cmod(q-r)*sin(angle-c q r p)⟩
  using ang-sin by fastforce
then have f4:⟨(p-r)*cnj(q-r)/cmod(q-r)^2 = cmod(p-r)*cmod(q-r)*cis(angle-c q r p)/cmod(q-r)^2⟩
apply(intro complex-eqI)
using f1 f2 by auto
then have f3:⟨(p-r)/(q-r) = cmod(p-r)*cmod(q-r)*cis(angle-c q r p)/cmod(q-r)^2⟩
by (simp add: ⟨(p - r) / (q - r) = (p - r) * cnj (q - r) / cor ((cmod (q - r)^2)))⟩
have ⟨cmod(p-r)*cmod(q-r)*cis(angle-c q r p)/cmod(q-r)^2 = cis (angle-c q r p)⟩
using ⟨p - r = cis (angle-c q r p) * (q - r)⟩ ⟨q ≠ r⟩ f3 by auto
then have ⟨cmod(p-r)*cmod(q-r)/cmod(q-r)^2 = 1⟩
by (metis f4 ⟨q ≠ r⟩ cmod-power2 complex-mod-cnj complex-mod-mult-cnj
complex-mult-cnj
eq-iff-diff-eq-0 nonzero-norm-divide norm-cis norm-eq-zero norm-mult power-not-zero)
then show ⟨cdist q r = cdist p r⟩
by (metis cdist-commute cdist-def eq-divide-eq-1 mult.commute power2-eq-square
real-divide-square-eq)
qed

```

```

lemma equilateral-imp-pi3:
assumes ⟨q≠r⟩ cdist q r = cdist p r cdist p r = cdist p q
shows |(angle-c q r p)| = pi/3 ∨ |(angle-c q r p)| = -pi/3
  (angle-c q r p) = (angle-c p q r) ∧ (angle-c q r p) = (angle-c r p q)
proof –
have f0:⟨q≠r ∧ r≠p ∧ p≠q⟩ using assms by(auto)
have ⟨angle-c q r p ≠ 0 ∧ angle-c q r p ≠ pi⟩
by (smt (verit, best) Re-complex-of-real angle-c-commute angle-c-commute-pi
assms(1) assms(2)
assms(3) cdist-commute cdist-def congruent-ctriangle-sss(20) cor-cmod-real
isoscele-iff-pr-cis-qr
minus-complex.simps(1) minus-complex.simps(2) minus-diff-eq mult-minus-right
norm-eq-zero right-minus-eq)
then have f1:⟨¬collinear q r p⟩
using collinear-angle f0 by blast
then have f1':⟨¬collinear p q r⟩
by (simp add: collinear-sym1 collinear-sym2)
have f12:⟨(angle-c q r p) = angle-c r p q ∨ (angle-c q r p) = - angle-c r p q⟩
apply(rule congruent-ctriangle-sss)
using f0 assms by(auto simp: norm-minus-commute)
have f4:⟨(angle-c q r p) = angle-c p q r ∨ (angle-c q r p) = - angle-c p q r⟩
apply(rule congruent-ctriangle-sss)
using f0 assms by(auto simp: norm-minus-commute)
have f5:⟨(angle-c q r p) ≠ pi⟩

```

```

using f1
by (metis angle-c-commute canon-ang-sin canon-ang-uminus-pi collinear-sin-neq-0
f1' pi-canonical
      sin-pi)
from f4 have ⟨ |(angle-c q r p) + (angle-c r p q) + (angle-c p q r)| = pi ⟩
using angle-sum-triangle-c[of p q r] f0 by auto
have ⟨ $\forall x \in \{-pi..0\}. 3*x \neq pi$ ⟩
by (metis atLeastAtMost-iff dual-order.trans mult-eq-0-iff mult-less-cancel-right2
numeral-le-one-iff
      pi-ge-zero pi-neq-zero semiring-norm(70) verit-comp-simplify1(3) verit-la-disequality)
also have the-x:⟨ $\exists!x \in \{-pi..pi\}. 3*x = pi$ ⟩
by(auto intro:exI[where x=⟨pi/3⟩])
moreover have ⟨(angle-c x y z) ∈ {−pi..pi}⟩ for x y z
using ang-c-in by auto
have ⟨ $x \in \{-3*pi..3*pi\} \wedge x - 2*pi = pi \implies (x = 3*pi)$ ⟩ for x k::real
by(auto simp:)
have ⟨ $x \in \{-3*pi..3*pi\} \wedge x - 2*k*pi = pi \wedge (k \geq 2 \vee k \leq -2) \implies False$ ⟩ for
x and k::int
proof −
assume h:⟨ $x \in \{-3*pi..3*pi\} \wedge x - 2*k*pi = pi \wedge (k \geq 2 \vee k \leq -2)$ ⟩
then have f0:⟨ $r = 2 \implies x - 2*r*pi \leq -pi$ ⟩ for r by auto
also have f1:⟨ $k > 2 \implies x - 2*k*pi < x - 2*2*pi$ ⟩
by(auto)
have f3:⟨ $x - 2*2*pi \leq -pi$ ⟩
using f0[of 2] by(auto simp:abs-real-def)
then have f2:⟨ $k > 2 \implies x - 2*k*pi < -pi$ ⟩
using f1 by argo
ultimately have f3:⟨ $k \geq 2 \implies x - 2*k*pi \leq -pi$ ⟩
apply(cases ⟨abs k=2⟩) using f0 by(auto)
have f0':⟨ $r = -2 \implies x - 2*r*pi > pi$ ⟩ for r ::int
using h by(auto)
have ⟨ $x + 2*2*pi > pi$ ⟩ using h by auto
also have ⟨ $k < -2 \implies x - 2*k*pi > x - 2*(-2)*pi$ ⟩
using f0'[of 2]
proof −
assume a1:  $k < -2$ 
have f2:  $pi < x - \text{real-of-int}(-4) * pi$ 
using f0' by force
have  $k + k \leq -4$ 
using a1 by force
then show ?thesis
using f2 by (metis (no-types) diff-left-mono h mult.commute mult-2-right
of-int-le-iff ordered-comm-semiring-class.comm-mult-left-mono pi-ge-zero verit-comp-simplify1(3))
qed
ultimately have f4:⟨ $k \leq -2 \implies x - 2*k*pi > pi$ ⟩
apply(cases ⟨ $k = -2$ ⟩) using h
⟨ $k < -2 \implies x - 2 * -2 * pi < x - \text{real-of-int}(2 * k) * pi$ ⟩ by auto
show False
using f3 f4 h by force

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qed
then have possible-k:<x∈{−3*pi<..3*pi} ∧ x−2*k*pi = pi ⟹ k=−1 ∨ k=1
∨ k=0> for k::int and x
  by force
have ang-3-in:<y∈{−3*pi<..3*pi} ∧ |y| = pi ⟷ y=pi ∨ y = 3*pi ∨ y=−pi>
for y
  proof(rule iffI)
    assume hyp:<y ∈ {− 3 * pi <.. 3 * pi} ∧ |y| = pi>
    have <∃ k. y − 2*k*pi = pi>
      by (metis add-diff-cancel-left' diff-add-cancel divide-self-if field-sum-of-halves
mult.commute mult-1 mult-2
      pi-neq-zero times-divide-eq-right)
    with hyp show <y = pi ∨ y=3*pi ∨ y=−pi>
      using hyp possible-k canon-ang-pi-3pi[of y] canon-ang-id canon-ang-uminus-pi

      by (smt (verit, del-insts) canon-ang-minus-3pi-minus-pi greaterThanAtMost-iff)
next
  assume <y = pi ∨ y = 3 * pi ∨ y = − pi>
  then show <y ∈ {− 3 * pi <.. 3 * pi} ∧ |y| = pi>
    by (auto simp:canon-ang-uminus-pi canon-ang-pi-3pi)
qed
ultimately have <|3*(angle-c q r p)| = pi ⟹ angle-c q r p = −pi/3 ∨ angle-c
q r p = pi/3>
proof –
  assume <|3*(angle-c q r p)| = pi>
  have <3*(angle-c q r p) ∈ {−3*pi<..3*pi}>
    using ang-c-in by auto
  then have f0:<3*(angle-c q r p) = −pi ∨ 3*(angle-c q r p) = pi ∨ 3*(angle-c
q r p) = 3*pi>
    using ang-3-in <|3*(angle-c q r p)| = pi> by blast
  then have <3*(angle-c q r p) = 3*pi ⟹ False>
    using collinear-sin-neq-pi by (simp add: f5)
  then show <angle-c q r p = −pi/3 ∨ angle-c q r p = pi/3>
    using f0 by auto
qed
have f2:<|angle-c x y z| = angle-c x y z> for x y z::complex
  using ang-vec-bounded angle-c-def canon-ang-id by fastforce
then have <|3*(angle-c q r p)| = pi>
  apply (cases <(angle-c q r p) = − angle-c r p q>)
  apply (simp add:f1 f2)
  apply (metis add.commute add.right-inverse add.right-neutral angle-sum-triangle-c
canon-ang-uminus-pi f0 f2 f4 f5)
  by (smt (verit) <|angle-c q r p + angle-c r p q + angle-c p q r| = pi>
    <angle-c q r p = angle-c p q r ∨ angle-c q r p = − angle-c p q r>
    <angle-c q r p = angle-c r p q ∨ angle-c q r p = − angle-c r p q>
    canon-ang-uminus-pi
    f2)
  then have <angle-c q r p = −pi/3 ∨ angle-c q r p = pi/3>
    using <|3 * angle-c q r p| = pi ⟹ angle-c q r p = − pi / 3 ∨ angle-c q r p

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```

= pi / 3) by auto
then show <|angle-c q r p| = pi / 3 ∨ |angle-c q r p| = -pi / 3>
using f2 by auto
then show <angle-c q r p = angle-c p q r ∧ angle-c q r p = angle-c r p q>
by (smt (verit, del-insts) <|angle-c q r p + angle-c r p q + angle-c p q r| = pi>
f12 collinear-sin-neq-0 collinear-sym1 f1' f2 f4 f5 sin-pi)
qed

lemma isosceles-triangle-converse:
assumes angle-c a b c = angle-c c a b ¬collinear a b c
shows dist a c = dist b c
by (metis assms(1) assms(2) cdist-def collinear-sin-neq-0 collinear-sym1 congruent-ctriangle-asa(8)
congruent-ctriangle-asa(9) dist-complex-def law-of-sines mult-cancel-left non-collinear-independant)

lemma pi3-imp-equilateral:
assumes <q≠r> <p≠q> <r≠p>
and <(angle-c q p r) = pi/3 ∨ (angle-c q p r) = -pi/3>
and <(angle-c q p r) = (angle-c r q p)>
and <(angle-c q p r) = (angle-c p r q)>
shows <cdist p r = cdist q r ∧ cdist p r = cdist p q>
proof(safe)
have f0:<¬ collinear q p r>
using assms(1–3) unfolding collinear-def
by (smt (verit, ccfv-SIG) arccos-one-half arcsin-one-half arcsin-plus-arccos
assms(4)
collinear-angle collinear-def divide-le-eq-1-pos divide-pos-pos minus-divide-left
pi-def pi-gt-zero pi-half)
then show <cdist p r = cdist q r>
using isosceles-triangle-converse[of q p r] assms
by (simp add: dist-complex-def norm-minus-commute)
show <cdist p r = cdist p q>
using f0 isosceles-triangle-converse[of p r q] assms
by (metis <cdist p r = cdist q r> cdist-def collinear-iff dist-commute dist-norm)
qed

lemma pi3-isoscele-imp-equilateral:
assumes <q≠r> <p≠q> cdist q r = cdist p r
and <|(angle-c q p r)| = pi/3 ∨ |(angle-c q p r)| = -pi/3>
shows <cdist p q = cdist r q>
proof(–)
have f:<angle-c p q r = (angle-c r p q)>
by (smt (verit, best) angle-c-commute angle-c-commute-pi assms(2) assms(3)
cdist-commute
collinear-angle isosceles-triangles orientation-respect-letter-order)
have f0:<r≠p>

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using assms(1) assms(3) by force
then have ⟨cdist q p = cdist r q⟩
by (smt (verit) congruent-ctriangle-sss(19) add-divide-distrib ang-vec-bounded
angle-c-commute angle-c-def angle-sum-triangle-c arccos-minus-one-half ar-
ccos-one-half arcsin-minus-one-half
arcsin-one-half arcsin-plus-arccos assms(2) assms(3) assms(4) canon-ang-id
canon-ang-minus-3pi-minus-pi
cdist-commute field-sum-of-halves law-of-sines mult-cancel-right pi3-imp-equilateral
sin-gt-zero sin-periodic-pi)
then show ⟨cdist p q = cdist r q⟩
using cdist-commute by force
qed

```

```

lemma pi3-isoscele-imp-equilateral':
assumes ⟨q ≠ r⟩ ⟨p ≠ q⟩ cdist q r = cdist p r
and |⟨angle-c q p r⟩| = pi/3 ∨ |⟨angle-c q p r⟩| = -pi/3
shows ⟨cdist p r = cdist p q⟩
by (metis assms(2) assms(3) assms(4) cdist-commute pi3-isoscele-imp-equilateral)

```

```

lemma equilateral-characterization:⟨q ≠ r ⟹ (cdist q r = cdist p r ∧ cdist p r =
cdist p q)⟩
↔ ((p - r) = cis(pi/3)*(q - r) ∨ (p - r) = cis(-pi/3)*(q - r))
proof(rule iffI)
assume h: ⟨q ≠ r⟩ ⟨cdist q r = cdist p r ∧ cdist p r = cdist p q⟩
then have f1:⟨p - r = cis(⟨angle-c q r p⟩) * (q - r) ∨ p - r = cis(-⟨angle-c
q r p⟩) * (q - r)⟩
using isoscele-iff-pr-cis-qr by blast
have f0:⟨q ≠ r ∧ r ≠ p ∧ p ≠ q⟩ using h by auto
have ⟨angle-c q r p = pi/3 ∨ angle-c q r p = -pi/3⟩
using equilateral-imp-pi3(1)[of q r p]
by (metis ang-vec-bounded angle-c-def canon-ang-id h(1) h(2))
then show ⟨p - r = cis(pi / 3) * (q - r) ∨ p - r = cis(-pi / 3) * (q - r)⟩
using f1
by (metis add.inverse-inverse minus-divide-left)
next
assume h:⟨q ≠ r⟩ ⟨p - r = cis(pi / 3) * (q - r) ∨ p - r = cis(-pi / 3) *
(q - r)⟩
have ⟨(p - r)/(q - r) = (p - r)*cnj(q - r)/cmod(q - r) ^ 2⟩
using complex-div-cnj by blast
have f1:⟨Re((p - r)*cnj(q - r)) = cmod(p - r)*cmod(q - r)*cos(⟨angle-c q r p⟩)⟩
using ang-cos by force
have f2: ⟨Im((p - r)*cnj(q - r)) = cmod(p - r)*cmod(q - r)*sin(⟨angle-c q r p⟩)⟩
using ang-sin h(1) by auto
then have ⟨(p - r)*cnj(q - r)/cmod(q - r) ^ 2 = cmod(p - r)*cmod(q - r)*cis(⟨angle-c
q r p⟩)/cmod(q - r) ^ 2⟩
apply(intro complex-eqI)

```

```

    using f1 f2 by auto
then have f3:( $(p - r) / (q - r) = \text{cmod}(p - r) * \text{cmod}(q - r) * \text{cis}(\text{angle-c } q \ r \ p) / \text{cmod}(q - r)^2$ )
    by (simp add:  $\langle (p - r) / (q - r) = (p - r) * \text{cnj}(q - r) / \text{cor}((\text{cmod}(q - r))^2)$ )
have  $\langle (p - r) / (q - r) = \text{cis}(\text{angle-c } q \ r \ p)$ 
apply(subst f3)
by (smt (z3)  $\langle (p - r) * \text{cnj}(q - r) / \text{cor}((\text{cmod}(q - r))^2) = \text{cor}(\text{cmod}(p - r) * \text{cmod}(q - r))$ 
     $* \text{cis}(\text{angle-c } q \ r \ p) / \text{cor}((\text{cmod}(q - r))^2)$ ) cis-divide cis-mult cis-neq-zero
cis-zero
cmod-cor-divide complex-mod-cnj eq-iff-diff-eq-0 f3 h(1) h(2) nonzero-mult-divide-mult-cancel-left
nonzero-mult-divide-mult-cancel-right norm-cis norm-mult numeral-One
of-real-1 of-real-divide
power2-less-0 times-divide-eq-left)
then have fpi: $\langle \text{angle-c } q \ r \ p = \pi/3 \vee \text{angle-c } q \ r \ p = -\pi/3 \rangle$ 
using h ang-c-in
by (smt (verit, best) add-divide-distrib ang-vec-bounded angle-c-def arccos-one-half
arcsin-one-half arcsin-plus-arccos cis-inj diff-eq-diff-eq diff-numeral-special(9)

divide-nonneg-pos field-sum-of-halves nonzero-divide-eq-eq)
then have ff: $\langle \text{cdist } p \ r = \text{cdist } q \ r \rangle$ 
by (metis  $\langle (p - r) / (q - r) = \text{cis}(\text{angle-c } q \ r \ p) \rangle$  h(1) isoscele-iff-pr-cis-qr
nonzero-divide-eq-eq right-minus-eq)
then have  $\langle \text{angle-c } p \ q \ r = \text{angle-c } q \ p \ r \vee \text{angle-c } p \ q \ r = -\text{angle-c } q \ p \ r \rangle$ 
by (metis congruent-ctriangle-sss(13) cdist-commute cdist-def norm-eq-zero
right-minus-eq)
then have f10: $\langle |\text{angle-c } p \ q \ r + \text{angle-c } q \ r \ p + \text{angle-c } r \ p \ q| = \pi \rangle$ 
by (metis  $\langle \text{angle-c } q \ r \ p = \pi/3 \vee \text{angle-c } q \ r \ p = -\pi/3 \rangle$  angle-c-neq0
angle-sum-triangle-c
divide-eq-0-iff h(1) neg-equal-0-iff-equal pi-neq-zero zero-neq-numeral)
then have  $\langle |\text{angle-c } q \ p \ r| = \pi/3 \vee |\text{angle-c } q \ p \ r| = -\pi/3 \rangle$ 
by (smt (verit)  $\langle \text{angle-c } p \ q \ r = \text{angle-c } q \ p \ r \vee \text{angle-c } p \ q \ r = -\text{angle-c } q \ p \ r \rangle$ 
add-divide-distrib ang-pos-pos ang-vec-bounded ang-vec-sym angle-c-def arc-
cos-minus-one-half
arccos-one-half arcsin-minus-one-half arcsin-one-half arcsin-plus-arccos
canon-ang-id
canon-ang-minus-3pi-minus-pi canon-ang-pi-3pi field-sum-of-halves fpi)
show  $\langle \text{cdist } q \ r = \text{cdist } p \ r \wedge \text{cdist } p \ r = \text{cdist } p \ q \rangle$ 
apply(rule conjI)
using ff apply(auto simp: norm-minus-commute)[1]
apply(rule pi3-isoscele-imp-equilateral')
using h(1) apply auto[1]
apply (metis  $\langle \text{angle-c } q \ r \ p = \pi/3 \vee \text{angle-c } q \ r \ p = -\pi/3 \rangle$ 
angle-c-neq0 divide-eq-0-iff neg-equal-0-iff-equal pi-neq-zero zero-neq-numeral)
using ff apply presburger
using  $\langle |\text{angle-c } q \ p \ r| = \pi/3 \vee |\text{angle-c } q \ p \ r| = -\pi/3 \rangle$  by blast

```

```

qed

lemmas equilateral-imp-prcisi3 = equilateral-characterization[THEN iffD1]

lemmas prcisi3-imp-equilateral = equilateral-characterization[THEN iffD2]

lemma equilateral-characterization-neg:
  fixes p q r::complex
  assumes h1:( $p \neq r$ )
  shows  $\langle cdist p r = cdist p q \wedge cdist p q = cdist q r \wedge angle-c q r p = -pi/3 \rangle$ 
     $\longleftrightarrow p + root3 * q + root3^2 * r = 0$ 
proof(rule iffI)
  assume h:( $cdist p r = cdist p q \wedge cdist p q = cdist q r \wedge angle-c q r p = -pi/3$ )
  then have  $\langle p - r = cis(-pi/3)*(q - r) \rangle$ 
    using h1 isoscele-iff-pr-cis-qr[of q r p]
    by force
  with h have  $\langle p - r + cis(2*pi/3)*(q - r) = 0 \rangle$ 
    apply(simp add:cis.code)
    apply(subst sin-120)
    apply(intro complex-eqI)
    using cos-120 cos-60 sin-120 sin-60
    by(auto simp add:field-simps)
  then have  $\langle p + root3*q - (root3+1)*r = 0 \rangle$  by(auto simp:field-simps root3-def)
  then show  $\langle p + root3 * q + root3^2 * r = 0 \rangle$ 
    by (metis (no-types, lifting) add.commute eq-iff-diff-eq-0 left-diff-distrib
      ring-class.ring-distrib(2) root3-eq-0)

next
  assume h:( $p + root3 * q + root3^2 * r = 0$ )
  then have  $\langle p + root3*q - (root3+1)*r = 0 \rangle$ 
    by (metis add.commute add-diff-cancel-left diff-numeral-special(9) left-diff-distrib
      ring-class.ring-distrib(2)
      root3-eq-0)
  then have  $\langle p - r + root3*(q - r) = 0 \rangle$ 
    by(auto simp:field-simps)
  have  $\langle -root3 = cis(-pi/3) \rangle$ 
    using cos-120 cos-60 sin-120 sin-60
    by (auto intro:complex-eqI simp: root3-def)
  then have f1':( $p - r = cis(-pi/3)*(q - r)$ )
    by (metis < $p - r + root3 * (q - r) = 0$ > eq-neg-iff-add-eq-0 mult-minus-left)
  then have f1:( $p - r = cis(pi/3)*(q - r) \vee p - r = cis(-pi/3)*(q - r)$ )
    by auto
  then have f2:( $cmod(p - r) = cmod(q - r)$ )
    by(auto simp:norm-mult)
  then have f3:( $dist p r = dist q r$ )
    by (simp add: dist-complex-def)
  have ang-dem:( $angle-c q r p = -pi/3$ )
  proof -
    have *:( $q \neq r$ )
      using f2 h1 by auto

```

```

then have **:( $(p - r) / (q - r) = cis(-pi/3)$ )
  using f1' by auto
have **:( $\text{angle-c } q \ r \ p = \text{Arg}((p - r) / (q - r))$ )
  by (metis ang-vec-def angle-c-def arg-div f1' h1 mult-zero-right right-minus-eq)
then have  $\langle \text{Arg}((p - r) / (q - r)) = -pi/3 \rangle$ 
  by (simp add:arg-cis **)
then show ?thesis using * ** ***
  using canon-ang-id pi-ge-two by force
qed
have  $\langle \text{cmod}(q - p) = \text{cmod}(r - q) \rangle$ 
  using f2
  by (metis cdist-def dist-eq-0-iff equilateral-characterization f1' f3)
then show  $\langle \text{cdist } p \ r = \text{cdist } p \ q \wedge \text{cdist } p \ q = \text{cdist } q \ r \wedge \text{angle-c } q \ r \ p = -pi/3 \rangle$ 

  using equilateral-characterization[THEN iffD2, of q r p, OF - f1]
  using f2 ang-dem by(auto simp: f2 field-simps intro!:)
qed

```

```

end
theory Complex-Axial-Symmetry

```

```

imports Complex-Angles Complex-Triangles

```

```

begin

```

6 Axial symmetry in complex field

In the following we define the axial symmetry and prove basics properties.

```

context
fixes z1 z2 :: complex and  $\alpha \beta :: \text{complex}$ 
assumes neq0:( $z1 \neq z2$ )
defines  $\langle \alpha \equiv (z1 - z2) / (\text{cnj } z1 - \text{cnj } z2) \rangle$ 
defines  $\langle \beta \equiv (z2 * \text{cnj } z1 - z1 * \text{cnj } z2) / (\text{cnj } z1 - \text{cnj } z2) \rangle$ 
begin

definition axial-symmetry::( $\text{complex} \Rightarrow \text{complex}$ ) where
   $\langle \text{axial-symmetry } z \equiv \text{cnj } z * (\text{cnj } z1 - \text{cnj } z2) / (\text{cnj } z1 - \text{cnj } z2) + (z2 * \text{cnj } z1 - z1 * \text{cnj } z2) / (\text{cnj } z1 - \text{cnj } z2) \rangle$ 

lemma norm-alpha-eq-1:( $\text{cmod}(\alpha) = 1$ )
  by (auto simp: neq0 alpha-def beta-def)

lemma z1-inv:( $\text{axial-symmetry } z1 = z1$ )
  unfolding axial-symmetry-def
  by (metis (mono-tags, lifting) add-diff-eq add-divide-distrib complex-cnj-cancel-iff
    diff-add-cancel eq-iff-diff-eq-0 mult.commute neq0 nonzero-eq-divide-eq right-diff-distrib')

```

```

lemma z2-inv: $\langle\text{axial-symmetry } z2 = z2\rangle$ 
  unfolding axial-symmetry-def
  by (smt (z3) add-diff-cancel-left add-diff-cancel-left' add-divide-distrib axial-symmetry-def
    complex-cnj-cancel-iff diff-add-cancel divide-divide-eq-right eq-divide-imp mult-commute-abs
    right-minus-eq z1-inv)

lemma cmod-axial: $\langle\text{cmod } (\text{axial-symmetry } z - \text{axial-symmetry } z') = \text{cmod } (\alpha * (\text{cnj } z - \text{cnj } z'))\rangle$ 
  unfolding axial-symmetry-def
  by (auto simp:  $\alpha$ -def  $\beta$ -def) (simp add: diff-divide-distrib mult.commute right-diff-distrib')

lemma cmod-axial-inv: $\langle\text{cmod } (\text{axial-symmetry } z - \text{axial-symmetry } z') = \text{cmod } (z - z')\rangle$ 
  by (metis cmod-axial complex-cnj-diff complex-mod-cnj mult.commute
    mult.right-neutral norm- $\alpha$ -eq-1 norm-mult)

lemma axial-symmetry-dist1: $\langle\text{cdist } z1 z = \text{cdist } z1 (\text{axial-symmetry } z)\rangle$ 
  by (metis cdist-def cmod-axial-inv z1-inv)

lemma axial-symmetry-dist2: $\langle\text{cdist } z2 z = \text{dist } z2 (\text{axial-symmetry } z)\rangle$ 
  by (metis cdist-def cmod-axial-inv dist-commute dist-norm z2-inv)

lemma  $\alpha\beta$ : $\langle\alpha * \text{cnj } \beta + \beta = 0\rangle$ 
  by (simp add: add-divide-eq-iff  $\alpha$ -def  $\beta$ -def)

lemma involution-symmetry: $\langle\text{axial-symmetry } (\text{axial-symmetry } z) = z\rangle$ 
proof -
  have  $\langle\alpha * \text{cnj}(\alpha * \text{cnj } z + \beta) + \beta = \alpha * \text{cnj } \alpha * z + \alpha * \text{cnj } \beta + \beta\rangle$ 
    by (simp add: ring-class.ring-distrib(1))
  then show ?thesis unfolding axial-symmetry-def
    by (smt (verit, best)  $\alpha\beta$  add.right-neutral cmod-eq-one group-cancel.add1
      mult.commute mult-cancel-right2 norm- $\alpha$ -eq-1 times-divide-eq-left  $\alpha$ -def
       $\beta$ -def)
  qed

lemma arg- $\alpha$ : $\langle\text{Arg } \alpha = |\text{Arg } (z1 - z2)|\rangle$ 
  unfolding  $\alpha$ -def  $\beta$ -def
  by (smt (verit, del-insts) arg-cnj-not-pi arg-div arg-pi-iff canon-ang-id canon-ang-pi-3pi
    complex-cnj-cancel-iff complex-cnj-diff eq-cnj-iff-real eq-iff-diff-eq-0 neq0 pi-ge-two)

lemma Arg-invol: $\langle\text{Arg } (\text{axial-symmetry } (\text{axial-symmetry } z)) = \text{Arg } z\rangle$ 
  by (simp add: involution-symmetry)

lemma angle-sum-symmetry: $\langle z \neq z1 \implies |\text{angle-c } z z1 z2 + \text{angle-c } z2 z1| \leq \pi \rangle$ 

```

```

 $z) \vdash \text{angle-c } z \text{ } z1 \text{ (axial-symmetry } z)$ 
proof –
  assume  $\langle z \neq z1 \rangle$ 
  have  $\langle z2 - z1 \neq 0 \rangle$  using  $\text{angle-c-sum neq0}$  by  $\text{auto}$ 
  have  $\langle z1 - (\text{axial-symmetry } z) \neq 0 \rangle$ 
    by (metis  $\langle z \neq z1 \rangle \text{ eq-iff-diff-eq-0 involution-symmetry } z1\text{-inv}$ )
  then show ?thesis
    using  $\text{angle-c-sum}$ 
    by (metis  $\langle z2 - z1 \neq 0 \rangle \text{ right-minus-eq } z1\text{-inv}$ )
qed

lemma  $\text{angle-symmetry-eq-imp}:$ 
assumes  $h: \langle z1 \neq z \rangle \langle z2 \neq z \rangle$ 
shows  $\langle \text{angle-c } z \text{ } z1 \text{ } z2 = - \text{angle-c (axial-symmetry } z) \text{ } z1 \text{ } z2 \vee \text{angle-c } z \text{ } z1 \text{ } z2 = \text{angle-c (axial-symmetry } z) \text{ } z1 \text{ } z2 \rangle$ 
by (metis (mono-tags, lifting)  $\text{axial-symmetry-dist1 cdist-def cmod-axial-inv}$ 
       $\text{congruent-ctriangle-sss(22)} \text{ } h(1) \text{ neq0 } z2\text{-inv}$ )

lemma  $\text{angle-symmetry}:$ 
assumes  $h: \langle z1 \neq z \rangle \langle z2 \neq z \rangle$ 
  and  $\langle \text{angle-c } z \text{ } z1 \text{ } z2 = \text{angle-c (axial-symmetry } z) \text{ } z1 \text{ } z2 \rangle$ 
shows  $\langle z = \text{axial-symmetry } z \rangle$ 
proof –
  have  $\langle \text{cmod } (z - z1) = \text{cmod } (\text{axial-symmetry } z - z1) \rangle$ 
    using  $\text{axial-symmetry-dist1}$  by  $\text{auto}$ 
  have  $\langle \text{cmod } (z - z2) = \text{cmod } (\text{axial-symmetry } z - z2) \rangle$ 
    by (metis  $\text{axial-symmetry-dist2 cdist-def dist-commute dist-complex-def}$ )
  have  $\langle z - z1 = \text{axial-symmetry } z - z1 \rangle$ 
    by (metis  $\langle \text{cmod } (z - z1) = \text{cmod } (\text{local.axial-symmetry } z - z1) \rangle \text{ ang-cos}$ 
       $\text{ang-sin assms(3)}$ 
       $\text{complex-cnj-cnj complex-eq-iff eq-iff-diff-eq-0 mult-cancel-left neq0}$ )
  then show ?thesis
    by  $\text{simp}$ 
qed

lemma  $\text{line-is-inv}: \langle z \in \text{line } z1 \text{ } z2 \wedge z \neq z2 \wedge z \neq z1 \implies z = \text{axial-symmetry } z \rangle$ 
proof –
  assume  $\langle z \in \text{line } z1 \text{ } z2 \wedge z \neq z2 \wedge z \neq z1 \rangle$ 
  then have  $\langle \text{angle-c } z \text{ } z1 \text{ } z2 = 0 \vee \text{angle-c } z \text{ } z1 \text{ } z2 = pi \rangle$ 
  unfolding  $\text{line-def}$  using  $\text{neq0 collinear-angle}$ 
  using  $\text{collinear-sym1 collinear-sym2'}$  by  $\text{blast}$ 
  then show ?thesis
    by (smt (verit)  $\langle z \in \text{line } z1 \text{ } z2 \wedge z \neq z2 \wedge z \neq z1 \rangle$ 
       $\text{angle-c-commute angle-c-commute-pi angle-symmetry angle-symmetry-eq-imp}$ )
qed

lemma  $\text{dist-inv}: \langle \text{cdist } a \text{ } b = \text{cdist } (\text{axial-symmetry } a) \text{ } (\text{axial-symmetry } b) \rangle$ 
by (simp add: cmod-axial-inv)

```

lemma *collinear-inv*: **assumes** $\langle \text{collinear } a \ b \ c \rangle$ **and** $\langle a \neq b \wedge b \neq c \wedge c \neq a \rangle$
shows $\langle \text{collinear } (\text{axial-symmetry } a) \ (\text{axial-symmetry } b) \ (\text{axial-symmetry } c) \rangle$
proof –
have $\langle \text{angle-c } a \ b \ c = pi \vee \text{angle-c } a \ b \ c = 0 \rangle$
using *assms(1)* *assms(2)* **collinear-angle** **by** *blast*
then have $\langle \text{angle-c } (\text{axial-symmetry } a) \ (\text{axial-symmetry } b) \ (\text{axial-symmetry } c) \rangle$
 $= \text{angle-c } a \ b \ c$
 $\quad \vee \text{angle-c } (\text{axial-symmetry } a) \ (\text{axial-symmetry } b) \ (\text{axial-symmetry } c) = -\text{angle-c } a \ b \ c$
by (*metis congruent-ctriangle-sss(24)* *assms(2)* *dist-inv minus-equation-iff*)
then show *?thesis*
using $\langle \text{angle-c } a \ b \ c = pi \vee \text{angle-c } a \ b \ c = 0 \rangle$ *collinear-sin-neq-0* *collinear-sym1*
by *fastforce*
qed

lemma *axial-symmetry-eq-line*: $\langle z \neq z1 \wedge z \neq z2 \implies z = \text{axial-symmetry } z \implies z \in \text{line } z1 \ z2 \rangle$
proof –
assume $\langle z \neq z1 \wedge z \neq z2 \rangle$ $\langle z = \text{axial-symmetry } z \rangle$
then have *g0*: $\langle z = \alpha * \text{cnj } z + \beta \rangle$ **unfolding** *axial-symmetry-def*
by (*simp add: mult.commute alpha-def beta-def*)
then have *g1*: $\langle \text{cnj } z = (z - \beta) * \text{cnj } \alpha \rangle$
by (*smt (verit, ccfv-threshold)* *add-diff-cancel* *complex-cnj-cnj* *complex-cnj-diff*
alpha-def *beta-def*
complex-cnj-divide *mult.commute neq0 nonzero-eq-divide-eq* *right-minus-eq*
times-divide-eq-left)
also have *g2*: $\langle z = (\text{cnj } z - \text{cnj } \beta) * \alpha \rangle$
by (*metis calculation complex-cnj-cnj complex-cnj-diff complex-cnj-mult*)
have $\langle \alpha / (z1 - z2) = 1 / (\text{cnj } z1 - \text{cnj } z2) \rangle$
by (*auto simp: alpha-def beta-def*)
have $\langle \text{Im } w = (w - \text{cnj } w) / (2*i) \rangle$ **for** *w*
using *Im-express-cnj* **by** *blast*
then have $\langle \text{Im } (z2 * \text{cnj } z1) = (z2 * \text{cnj } z1 - \text{cnj } (z2 * \text{cnj } z1)) / (2*i) \rangle$
by *presburger*
then have $\langle \text{Im } (z2 * \text{cnj } z1) * 2*i = (z2 * \text{cnj } z1 - \text{cnj } z2 * z1) \rangle$
by (*metis complex-cnj-cnj complex-cnj-mult complex-diff-cnj mult.commute*)
have *f0*: $\langle (\text{cnj } z1 - \text{cnj } z2) * (z1 - z2) = \text{cmod } z1^2 + \text{cmod } z2^2 - \text{cnj } z1 * z2 - \text{cnj } z2 * z1 \rangle$
by (*smt (verit)* *cancel-ab-semigroup-add-class.diff-right-commute* *complex-mult-cnj-cmod*
diff-diff-eq2 *left-diff-distrib* *mult.assoc* *mult.commute* *norm-minus-commute*
norm-mult of-real-add *of-real-eq-iff*)
have *f1*: $\langle \text{cmod } z1^2 + \text{cmod } z2^2 - \text{cnj } z1 * z2 - \text{cnj } z2 * z1 = \text{cmod } z1^2 +$
 $\text{cmod } z2^2 - 2 * \text{Re } z1 * \text{Re } z2 - 2 * \text{Im } z1 * \text{Im } z2 \rangle$
by (*auto simp: field-simps intro: complex-eqI*)
have $\langle \beta = 2*i * \text{Im } (z2 * \text{cnj } z1) / (\text{cnj } z1 - \text{cnj } z2) \rangle$
using $\langle \text{cor } (\text{Im } (z2 * \text{cnj } z1)) = (z2 * \text{cnj } z1 - \text{cnj } (z2 * \text{cnj } z1)) / (2 * i) \rangle$
unfolding *alpha-def* *beta-def* **by** *auto*
then have *f2*: $\langle \beta / (z1 - z2) = 2*i * \text{Im } (z2 * \text{cnj } z1) / ((\text{cnj } z1 - \text{cnj } z2) * (z1 - z2)) \rangle$

```

    by simp
  then have < $\beta/(z_1 - z_2) = 2*i*Im(z_2*cnj z_1)/(cmod z_1^2 + cmod z_2^2 -$ 
 $2*Re z_1*Re z_2 - 2*Im z_1*Im z_2)$ >
    using f0 f1 by presburger
  then have <is-real ((z - z1)/(z1 - z2))>
    unfolding axial-symmetry-def using g2
  by (smt (verit, del-insts) add-diff-cancel-right axial-symmetry-def complex-cnj-cnj
      complex-cnj-diff complex-cnj-divide diff-divide-distrib divide-divide-eq-right
      eq-cnj-iff-real
      g0 g2 mult.commute times-divide-eq-right z1-inv z2-inv α-def β-def)
  then show <z ∈ line z1 z2>
    unfolding line-def collinear-def
  by (metis (mono-tags, lifting) Im-i-times Re-i-times cnj.simps(2) complex-i-mult-minus
      eq-cnj-iff-real mem-Collect-eq minus-diff-eq minus-divide-right)
qed

lemma angle-symmetry-eq:
  assumes h:<z1 ≠ z> <z2 ≠ z> <z ∉ line z1 z2>
  shows <angle-c z z1 z2 = - angle-c (axial-symmetry z) z1 z2>
proof -
  have f0:<angle-c z z1 z2 = - angle-c (axial-symmetry z) z1 z2 ∨
    angle-c z z1 z2 = angle-c (axial-symmetry z) z1 z2>
    using angle-symmetry-eq-imp h(1) h(2) by blast
  have f1:<angle-c z z2 z1 = - angle-c (axial-symmetry z) z2 z1 ∨
    angle-c z z2 z1 = angle-c (axial-symmetry z) z2 z1>
    by (metis axial-symmetry-dist1 congruent-ctriangle(sss)(24) dist-inv h(1) neq0
        z2-inv)
  show ?thesis
    using angle-symmetry axial-symmetry-eq-line f0 h(1) h(2) h(3) by presburger
qed

end
end
theory Morley

```

imports Complex-Axial-Symmetry

begin

7 Rotations

```

locale complex-rotation =
  fixes A::complex and θ::real
begin

```

definition < $r z = A + (z - A)*cis(\vartheta)$ >

```

lemma cmod-inv-rotation:⟨cmod (z-A) = cmod (r z - A)⟩
  unfolding r-def
  by (simp add: norm-mult)

lemma inner-ang:⟨cos (⟨z1 z2) * (cmod z1 * cmod z2) = Re (innerprod z1 z2)⟩
proof -
  have ⟨Re (innerprod z1 z2) = Re (scalprod z1 z2)⟩
    unfolding innerprod-def by (auto)
  then show ?thesis
    by (metis cos-cmod-scalprod mult.commute)
qed

lemma ang-eq-cos-theta:⟨z ≠ A ⟹ cos (angle-c z A (r z)) = cos (θ)⟩
proof -
  assume ⟨z ≠ A⟩
  then have ⟨innerprod (z-A) ((r z - A)) = (z-A)*cis(θ)*cnj(z - A)⟩
    unfolding innerprod-def r-def by auto
  then have f0:⟨innerprod (z-A) ((r z - A)) = cmod (z-A)^2*cis(θ)⟩
    by (metis ab-semigroup-mult-class.mult-ac(1) complex-mult-cnj-cmod mult.commute)
  then have ⟨cos (angle-c z A (r z)) * cmod (z-A) * cmod (z - A) = Re (innerprod
    (z-A) ((r z - A)))⟩
    unfolding angle-c-def
    by (metis inner-ang mult.assoc cmod-inv-rotation)
  then have ⟨cos (angle-c z A (r z)) = Re (cis θ)⟩
    by (simp add: ⟨z ≠ A⟩ f0 power2-eq-square)
  then show ?thesis
    by (auto simp:cis.code)
qed

lemma cdist-dist:⟨cdist = dist⟩
  using cdist-commute dist-complex-def by fastforce

lemma ang-eq-theta:assumes h:⟨z ≠ A⟩ shows ⟨angle-c z A (r z) = |θ|⟩
proof(cases ⟨angle-c z A (r z) = |- θ|⟩)
  case True
  then have ⟨r z = A + (z-A)*cis(-θ)⟩
    by (metis add-diff-cancel-left' arg-cis cdist-def cis-cmod cis-neq-zero cmod-inv-rotation
      isoscele-iff-pr-cis-qr mult.left-commute mult.right-neutral mult-cancel-left
      norm-cis
      norm-minus-commute of-real-1 r-def right-minus-eq)
  then show ?thesis
    by (smt (verit, del-insts) True add-diff-cancel-left' angle-c-neq0 arg-cis arg-mult-eq
      canon-ang-cos canon-ang-sin cis.code cis-cnj diff-add-cancel divisors-zero
      r-def)
  next
  case False
  have ⟨cos (angle-c z A (r z)) = cos (θ)⟩
    using ang-eq-cos-theta h by auto

```

```

then show ?thesis
  by (smt (verit, ccfv-threshold) cmod-inv-rotation r-def add-diff-cancel-left' angle-c-neq0
    angle-c-sum arg-cis canon-ang-sin cdist-def cis.code cis-neq-zero divisors-zero
    eq-iff-diff-eq-0 h
    isoscele-iff-pr-cis-qr nonzero-mult-div-cancel-left nonzero-mult-div-cancel-right
    norm-minus-commute)
qed

lemma inj-r:<inj r>
  unfolding inj-on-def by(auto simp:r-def)

lemma img-eqI:<cdist A z1 = cdist A z2 ∧ angle-c z1 A z2 = θ ⟹ z2 = r z1>
  apply(cases ‹z1 = A ∨ z2 = A›)
  using r-def add-minus-cancel isoscele-iff-pr-cis-qr apply force
  unfolding r-def
  by (metis add.commute cdist-commute diff-add-cancel isoscele-iff-pr-cis-qr mult.commute)

lemma r-id-iff:<|θ| = 0 ⟺ r = id>
proof –
  obtain cc :: (complex ⇒ complex) ⇒ complex
  and cca :: (complex ⇒ complex) ⇒ complex where
    f1: ∀f. (id = f ∨ f (cc f) ≠ cc f) ∧ ((∀c. f c = c) ∨ id ≠ f)
    by (metis (no-types) eq-id-iff)
  have f2: ∀c ca ra. ca + (c - ca) * Complex (cos ra) (sin ra) = complex-rotation.r ca ra c
    by (simp add: complex-rotation.r-def cis.code)
  then have ∀c ca. complex-rotation.r c 0 ca = ca
    by (metis (no-types) add-diff-cancel-left' cos-zero diff-diff-eq2
      lambda-one mult.commute one-complex.code sin-zero)
  then show ?thesis
    using f2 f1
    by (metis (lifting, full-types) ang-eq-theta angle-c-neq0
      canon-ang-cos canon-ang-sin complex-i-not-zero)
qed

end

lemma axial-symmetry-eq:<axial-symmetry B C P = axial-symmetry C B P> if
  ‹C ≠ B› for C B P
  unfolding axial-symmetry-def[OF that] axial-symmetry-def[OF that[symmetric]]]

by (metis (no-types, lifting) complex-cnj-cancel-iff eq-iff-diff-eq-0
  minus-diff-eq nonzero-minus-divide-divide times-divide-eq-right)

lemma img-r-sym:
  assumes h:<z1 ≠ z2› ‹z ∉ line z1 z2›
  shows ‹axial-symmetry z1 z2 z = complex-rotation.r z1 (|2*angle-c z z1 z2|) z›

```

```

proof -
  interpret complex-rotation z1 |2*angle-c z z1 z2| .
  let ?z = <axial-symmetry z1 z2 z>
  from h have <z1 ≠ z> <z2 ≠ z>
    unfolding line-def Elementary-Complex-Geometry.collinear-def by(auto)
  then have <angle-c ?z z1 z2 = -angle-c z z1 z2>
    using angle-symmetry-eq h(1) h(2) by force
  then have <angle-c z z1 ?z = |2*angle-c z z1 z2|>
    by (metis <z1 ≠ z> add.inverse-inverse angle-c-commute angle-c-commute-pi
      angle-sum-symmetry h(1) mult-2-right mult-commute-abs)
  have < cdist z1 z = cdist z1 (axial-symmetry z1 z2 z)>
    using axial-symmetry-dist1 h(1) by blast
  then show ?thesis
    apply(intro img-eqI)
    by (metis <angle-c z z1 (axial-symmetry z1 z2 z) = |2 * angle-c z z1 z2|>)
qed

lemma img-r-sym':
  assumes h:<z1 ≠ z2> <znotinline z1 z2>
  shows <axial-symmetry z1 z2 z = complex-rotation.r z1 (|-2*angle-c z2 z1 z|)>
  by (metis angle-c-commute angle-c-neq0 axial-symmetry-eq-line
    cancel-comm-monoid-add-class.diff-cancel complex-rotation.img-eqI
    complex-rotation.r-def h(1,2) img-r-sym mult-eq-0-iff mult-minus-left
    mult-minus-right pi-neq-zero two-pi-canonical)

lemma equality-for-pqr:
  assumes 1:<(a2::complex)*a3 ≠ 1> and 2:<∀z. h z = a3*z + b3> and 3:<∀z. g
  z = a2*z + b2> and 4:<g (h z) = z>
  shows <z = (a2*b3 + b2)/(1-a2*a3)>
proof -
  have f21:<g (h z) = a2*a3*z + a2*b3 + b2>
    using assms by(auto simp:2 3field-simps)
  then have <g (h z) = a2*a3*z + a2*b3 + b2 ↔ z*(1-a2*a3) = a2*b3 +
  b2>
    by(auto simp:field-simps 4)
  then have <a2*a3 ≠ 1 ⇒ z = (a2*b3 + b2)/(1-a2*a3)>
    using f21 by(auto simp:field-simps)
  then show ?thesis using 1 by auto
qed

lemma equality-for-comp:
  assumes 2:<∀z. h z = (a3::complex)*z + b3> and 3:<∀z. g z = a2*z + b2>
  and 4:<∀z. f z = a1*z + b1>
  shows <((f o f o f) o (g o g o g) o (h o h o h)) z = (a1*a2*a3)^3*z + (a1^2+a1+1)*b1
  + a1^3*(a2^2+a2+1)*b2
  + a1^3*a2^3*(a3^2+a3+1)*b3>
  using assms unfolding comp-def by(auto simp:fun-eq-iff power2-eq-square power3-eq-cube
  field-simps)

```

```

lemma eq-translation-id:
  assumes ⟨ $h = \text{complex-rotation.r } A \ 0$ ⟩ ⟨ $h B = B$ ⟩
  shows ⟨ $h = id$ ⟩
  using assms(1) complex-rotation.r-id-iff by auto

lemma r-eqI:
  assumes ⟨ $A = B$ ⟩ ⟨ $\vartheta_1 = \vartheta_2$ ⟩
  shows ⟨ $r A \vartheta_1 = r B \vartheta_2$ ⟩
  using assms(1) assms(2) by force

lemma r-eqI':
  assumes ⟨ $A = B$ ⟩ ⟨ $\vartheta_1 = \vartheta_2$ ⟩
  shows ⟨ $r A \vartheta_1 z = r B \vartheta_2 z$ ⟩
  using assms(1) assms(2) by force

lemma composed-rotations-same-center:
  shows ⟨ $(\text{complex-rotation.r } A \vartheta_1 \circ \text{complex-rotation.r } A \vartheta_2) = \text{complex-rotation.r } A (\vartheta_1 + \vartheta_2)$ ⟩
  unfolding complex-rotation.r-def by (auto simp: fun-eq-iff cis-mult add-ac)

lemma composed-rotations:
  assumes  $h : |\vartheta_1 + \vartheta_2| \neq 0$ 
  shows ⟨ $(\text{complex-rotation.r } A \vartheta_1 \circ \text{complex-rotation.r } B \vartheta_2) =$ 
     $\text{complex-rotation.r } ((A*(1 - cis \vartheta_1) + B*cis \vartheta_1*(1 - cis \vartheta_2))/(1 - cis(\vartheta_1 + \vartheta_2))) (\vartheta_1 + \vartheta_2)$ ⟩
  proof –
    have ⟨ $cis(\vartheta_1 + \vartheta_2) \neq 1$ ⟩
      by (metis arg-cis assms cis-zero zero-canonical)
    with  $h$  have ⟨ $(\text{complex-rotation.r } A \vartheta_1 \circ \text{complex-rotation.r } B \vartheta_2)$ 
       $((A*(1 - cis \vartheta_1) + B*cis \vartheta_1*(1 - cis \vartheta_2))/(1 - cis(\vartheta_1 + \vartheta_2)))$ 
       $= (A*(1 - cis \vartheta_1) + B*cis \vartheta_1*(1 - cis \vartheta_2))/(1 - cis(\vartheta_1 + \vartheta_2))$ ⟩
    unfolding complex-rotation.r-def using assms
    by(auto simp:cis-mult field-simps intro!:)
  with  $h$  show ?thesis
    apply(cases ⟨ $\vartheta_1 = 0 \vee \vartheta_2 = 0$ ⟩)
    unfolding complex-rotation.r-def using assms
    by(auto simp:cis-mult field-simps fun-eq-iff diff-divide-distrib add-divide-distrib
      intro!:)
  qed

lemma composed-rotation-is-trans:
  assumes ⟨ $|\vartheta_1 + \vartheta_2| = 0$ ⟩
  shows ⟨ $(\text{complex-rotation.r } A \vartheta_1 \circ \text{complex-rotation.r } B \vartheta_2) z = z + (B -$ 
     $A)*(cis(\vartheta_1) - 1)$ ⟩
  using assms
  by (auto simp:complex-rotation.r-def add-divide-distrib diff-divide-distrib field-simps)

```

(metis add.commute canon-ang-cos canon-ang-sin cis.code cis-mult cis-zero lambda-one mult.commute)

8 Morley's theorem

We begin by proving the Morley's theorem in the case where angles are positives then using the congruence between two triangles with the same angles only not of the same sign we prove Morley's theorem when angles are negatives.

We then proceed to conclude because in a triangle either angles are all negatives or all the angles are positives depending on orientation.

theorem *Morley-pos*:

```

assumes ¬collinear A B C
  ⟨angle-c A B R = angle-c A B C / 3⟩ (is ⟨?abr = ?abc⟩)
  angle-c B A R = angle-c B A C / 3 (is ⟨?bar = ?α⟩)
  angle-c B C P = angle-c B C A / 3 (is ⟨?bcp = ?bca⟩)
  angle-c C B P = angle-c C B A / 3 (is ⟨?cbp = ?β⟩)
  angle-c C A Q = angle-c C A B / 3 (is ⟨?caq = ?cab⟩)
  angle-c A C Q = angle-c A C B / 3 (is ⟨?acq = ?γ⟩)
  and hhh:⟨|angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3| = pi/3⟩
shows ⟨cdist R P = cdist P Q ∧ cdist Q R = cdist P Q⟩
proof –
  have bundle-line:⟨A ∉ line B C⟩ ⟨B ∉ line A C⟩ ⟨C ∉ line A B⟩ ⟨A ≠ B⟩ ⟨B ≠ C⟩
  ⟨C ≠ A⟩
  using assms(1) non-collinear-independant by (auto simp:collinear-sym1 collinear-sym2
  line-def)
  {fix A B C γ
    assume ABC-nline:⟨A ∉ line B C⟩
    and eq-3c:⟨angle-c A C B = 3*γ⟩
    then have neq-PI:⟨abs γ < pi/3⟩
    proof –
      have ⟨angle-c A C B ≠ pi⟩
      using ABC-nline(1) unfolding line-def
      by (metis angle-c-commute-pi collinear-iff mem-Collect-eq non-collinear-independant)
      then have ⟨angle-c A C B ∈ {−pi < .. < pi}⟩
      using ang-c-in less-eq-real-def by auto
      then show ⟨abs γ < pi/3⟩
      using eq-3c by force
    qed}note ang-inf-pi3=this
  have ⟨|angle-c B A C + angle-c C B A + angle-c A C B| = pi⟩
  by (metis collinear-def add.commute
        angle-sum-triangle-c assms(1) collinear-sym1 collinear-sym2')
  have eq-pi:⟨|3*?α + 3*?β + 3*?γ| = pi⟩
  using ⟨|angle-c B A C + angle-c C B A + angle-c A C B| = pi⟩ by force
  then have neq-pi:⟨|?β| ≠ pi ∧ |?γ| ≠ pi ∧ |?α| ≠ pi⟩
  by (smt (verit) ang-neg-neg angle-c-commute angle-c-neq0 assms

```

canon-ang-sin collinear-angle collinear-sin-neq-0 divide-eq-0-iff divide-nonneg-pos

```

minus-divide-left pi-neq-zero sin-pi sin-pi-minus zero-canonical)
moreover have <math>\langle 3*\alpha \neq 0 \wedge 3*\beta \neq 0 \wedge 3*\gamma \neq 0 \rangle
  using bundle-line collinear-sin-neq-0 line-def angle-c-commute assms(1) bun-
  dle-line(4)
    bundle-line(5) bundle-line(6) collinear-iff assms(1) collinear-sin-neq-0
    by (metis collinear-sym2' divide-eq-eq-numeral1(1) mem-Collect-eq mult.commute
zero-neq-numeral)
  ultimately have neq-0:<math>\langle |\beta| \neq 0 \wedge |\gamma| \neq 0 \wedge |\alpha| \neq 0 \rangle
    by (metis ang-vec-def angle-c-def arg-cis assms(3) assms(5) assms(7) canon-ang-arg
mult-zero-right)

```

interpret rot1: complex-rotation A $2*\alpha$.

interpret rot2: complex-rotation B $2*\beta$.

interpret rot3: complex-rotation C $2*\gamma$.

```

let ?f=<rot1.r>
have f0:<rot1.r A = A>
  unfolding rot1.r-def by auto
have f1:<rot2.r B = B>
  unfolding rot2.r-def by auto
have f2:<rot3.r C = C>
  unfolding rot3.r-def by auto

have <math>\langle cmod(rot3.r z - C) = cmod(z - C) \rangle \text{ for } z
  using rot3.cmod-inv-rotation by presburger
have f2:<math>\langle B \neq C \rangle
  by (metis collinear-def assms(1) collinear-sym1 collinear-sym2)
have f5:<math>\langle angle-c C B P = \beta \rangle \langle angle-c P B C = -\beta \rangle
  by (auto simp add: assms(5) angle-c-commute)
    (metis angle-c-commute angle-c-commute-pi assms(5) neq-pi pi-canonical)
then have f3:<math>\langle P \notin line C B \rangle
  by (smt (verit, ccfv-SIG) neq-pi angle-c-commute angle-c-neq0 assms(4)
collinear-angle divide-eq-0-iff line-def mem-Collect-eq pi-canonical zero-canonical)
then have f3':<math>\langle P \notin line B C \rangle
  using collinear-sym2' line-def by blast
then have f4:<math>\langle P \neq C \wedge P \neq B \rangle
  by (metis collinear-def collinear-sym1 collinear-sym2 line-def mem-Collect-eq)
then have <math>\langle angle-c P B C = - angle-c (axial-symmetry B C P) B C \rangle
  using angle-symmetry-eq[OF f2 - - f3] by auto
then have <math>\langle angle-c (axial-symmetry B C P) B C = \beta \rangle
  using f5(2) by fastforce

have f13:<math>\langle angle-c P C B = \gamma \rangle
  by (smt (verit, ccfv-threshold) angle-c-commute angle-c-commute-pi assms(1)
assms(4) assms(7)
collinear-sin-neq-0 nonzero-minus-divide-divide nonzero-minus-divide-right
sin-pi)

```

```

let ?P' = ⟨(axial-symmetry B C P)⟩

have ⟨angle-c (axial-symmetry B C P) B P = |2*?β|⟩
  by (metis ⟨P ≠ C ∧ P ≠ B⟩ ⟨angle-c (axial-symmetry B C P) B C = angle-c C B A / 3⟩
    angle-sum-symmetry f2 f5(1) involution-symmetry mult.commute mult-2-right
    z1-inv)
have ⟨cdist B (P) = cdist B ?P'⟩
  by(auto)(metis cmod-axial-inv f2 z1-inv)
then have ⟨cdist B (rot2.r P) = cdist B ?P'⟩
  unfolding cdist-def
  using rot2.cmod-inv-rotation by presburger
have f16:⟨rot2.r ?P' = P⟩
  by (metis ⟨angle-c (axial-symmetry B C P) B P = |2 * (angle-c C B A / 3)|⟩
    ⟨cdist B P = cdist B (axial-symmetry B C P)⟩ canon-ang-cos
    canon-ang-sin cis.code complex-rotation.img-eqI complex-rotation.r-def)
have f10:⟨angle-c P C B = |?γ|⟩
  by (metis ⟨angle-c P C B = angle-c A C B / 3⟩ ang-vec-def angle-c-def arg-cis
    canon-ang-arg)

from f10 have f11:⟨angle-c B C ?P' = |?γ|⟩
  by (metis axial-symmetry-eq ⟨P ≠ C ∧ P ≠ B⟩ ⟨angle-c P C B = angle-c A C
    B / 3⟩ angle-c-commute
    angle-symmetry angle-symmetry-eq-imp axial-symmetry-eq-line canon-ang-uminus-pi
    f2 f3 pi-canonical)
then have f12:⟨angle-c P C ?P' = |2*?γ|⟩
  by (metis axial-symmetry-eq f13 angle-sum-symmetry f10 f2 f4 mult.commute
    mult-2-right)
then have f15:⟨rot3.r P = ?P'⟩
  by (metis axial-symmetry-eq canon-ang-cos canon-ang-sin cis.code f13 f2 f3
    img-r-sym complex-rotation.r-def)
have P-inv:⟨rot3.r (rot3.r P) = P⟩
  using f16 f15 by presburger
let ?Q' = ⟨axial-symmetry A C Q⟩
have ⟨angle-c C A Q = -?α⟩
  by (metis angle-c-commute assms(1) assms(6) collinear-sin-neq-0
    collinear-sym1 collinear-sym2' minus-divide-left sin-pi)
then have ⟨angle-c Q A C = ?α⟩
  by (metis add.inverse-inverse angle-c-commute canon-ang-uminus-pi neq-pi
    pi-canonical)
have ⟨A ≠ C ∧ Q ∉ line A C⟩
  by (metis ⟨angle-c Q A C = angle-c B A C / 3⟩ angle-c-neq0 assms(7)
    collinear-angle div-0
    f10 f13 line-def mem-Collect-eq neq-0 neq-pi zero-canonical)
then have f17:⟨rot1.r Q = ?Q'⟩
  using img-r-sym[of A C Q]
  using ⟨angle-c Q A C = angle-c B A C / 3⟩ canon-ang-cos canon-ang-sin
    cis.code complex-rotation.r-def
  by presburger

```

then have $\langle \text{angle-c } ?Q' C A = ?\gamma \rangle$
by (smt (verit, ccfv-SIG) $\langle A \neq C \wedge Q \notin \text{line } A C \rangle$ $\langle \text{angle-c } Q A C = \text{angle-c } B A C / 3 \rangle$
 $\text{complex-rotation.ang-eq-theta angle-c-commute angle-symmetry angle-symmetry-eq-imp}$
 $\text{assms}(7)$
 $\text{axial-symmetry-eq axial-symmetry-eq-line f10 f13 complex-rotation.img-eqI}$
 $\text{neq-0 neq-pi})$
then have $\langle \text{angle-c } A C Q = ?\gamma \rangle$
using $\text{assms}(7)$ **by** blast
then have $\langle \text{angle-c } ?Q' C Q = |2 * ?\gamma| \rangle$
by (metis $\langle A \neq C \wedge Q \notin \text{line } A C \rangle$ $\langle \text{angle-c } (\text{axial-symmetry } A C Q) C A = \text{angle-c } A C B / 3 \rangle$
 $\langle \text{angle-c } Q A C = \text{angle-c } B A C / 3 \rangle$ $\text{angle-c-neq0 angle-sum-symmetry}$
 angle-symmetry-eq
 $\text{axial-symmetry-eq involution-symmetry mult-2-right mult-commute-abs neg-equal-0-iff-equal}$
 $\text{neq-0 zero-canonical})$
then have $f18:\langle \text{rot3.r } ?Q' = Q \rangle$
using $\text{img-r-sym}[\text{of } C A ?Q']$
by (metis $\langle A \neq C \wedge Q \notin \text{line } A C \rangle$ $\text{axial-symmetry-dist1}$
 $\text{axial-symmetry-eq canon-ang-cos canon-ang-sin cis.code}$
 $\text{complex-rotation.img-eqI complex-rotation.r-def})$
have $Q\text{-inv}:\langle \text{rot3.r } (\text{rot1.r } Q) = Q \rangle$
using $f17 f18$ **by** presburger
let $?R'=\langle \text{axial-symmetry } A B R \rangle$
have $\langle \text{angle-c } B A R = ?\alpha \rangle$
by (simp add: $\text{assms}(3)$)
then have $\langle \text{angle-c } R A B = -?\alpha \rangle$
by (metis angle-c-commute canon-ang-uminus-pi neq-pi pi-canonical)
have $\langle \text{angle-c } R B A = ?\beta \rangle$
by (metis $\langle \text{angle-c } (\text{axial-symmetry } B C P) B C = \text{angle-c } C B A / 3 \rangle$ angle-c-commute
 $\text{angle-c-commute-pi assms}(1) \text{ assms}(2) \text{ collinear-sin-neq-0 collinear-sym1}$
 $\text{minus-divide-left sin-pi})$
have $\langle A \neq B \wedge R \notin \text{line } A B \rangle$
by (metis (no-types, opaque-lifting) $\langle \text{angle-c } Q A C = \text{angle-c } B A C / 3 \rangle$
 $\langle \text{angle-c } R A B = -(\text{angle-c } B A C / 3) \rangle$ angle-c-commute angle-c-commute-pi angle-c-neq0 assms(2)
 collinear-angle
 $\text{collinear-sym2}' \text{ diff-zero divide-eq-0-iff line-def mem-Collect-eq minus-diff-eq}$
 neg-equal-zero
 $\text{neq-0 zero-canonical})$
then have $f17:\langle \text{rot2.r } R = ?R' \rangle$
using $\text{img-r-sym}[\text{of } B A R]$
using $\langle \text{angle-c } R B A = \text{angle-c } C B A / 3 \rangle$ axial-symmetry-eq
 $\text{cis.code collinear-sym2 line-def complex-rotation.r-def by auto}$
then have $\langle \text{angle-c } ?R' A B = ?\alpha \rangle$
by (metis $\langle A \neq B \wedge R \notin \text{line } A B \rangle$ $\langle \text{angle-c } R A B = -(\text{angle-c } B A C / 3) \rangle$
 $\langle \text{angle-c } R B A = \text{angle-c } C B A / 3 \rangle$ add.inverse-inverse add.inverse-add.inverse-neutral
 angle-c-neq0

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angle-symmetry-eq neq-0 zero-canonical)
then have f18:<rot1.r ?R' = R>
  using img-r-sym[of A B ?R]
  by (metis ‹A ≠ B ∧ R ∉ line A B› axial-symmetry-eq canon-ang-cos canon-ang-sin
cis.code
    involution-symmetry line-is-inv complex-rotation.r-def z2-inv)
have R-inv:<rot1.r (rot2.r R) = R>
  using f17 f18 by presburger
let ?a1 = ‹cis(2*?α)› let ?b1 = ‹A*(1 - ?a1)› let ?a2 = ‹cis(2*?β)› let ?b2 = ‹B*(1 - ?a2)›
let ?a3 = ‹cis(2*?γ)› let ?b3 = ‹C*(1 - ?a3)›
have possi-abc:<|?α + ?β + ?γ| = pi/3 ∨ |?α + ?β + ?γ| = -pi/3 ∨ |?α + ?β + ?γ| =
pi›
proof -
  have <|?α + ?β + ?γ| ∈ {-pi..pi}>
    by (simp add: canon-ang(1) canon-ang(2))
  then have <3 * |?α + ?β + ?γ| ∈ {-3*pi..3*pi}>
    by (auto)
  then have <||?α + ?β + ?γ| + |?α + ?β + ?γ| + |?α + ?β + ?γ|| = |?α + ?β + ?γ + ?α + ?β + ?γ + ?α + ?β + ?γ|>
    by (smt (verit, ccfv-SIG) add.commute arg-cis canon-ang-arg
      canon-ang-sum distrib-left is-num-normalize(1) mult-2)
  then have <|3 * |?α + ?β + ?γ|| = pi>
    using eq-pi by argo
  then have <3 * |?α + ?β + ?γ| = 3*pi ∨ 3 * |?α + ?β + ?γ| = -pi ∨ 3 * |?α + ?β + ?γ| = pi>
    by (smt (verit, del-insts) canon-ang(1) canon-ang(2) canon-ang-id
      canon-ang-minus-3pi-minus-pi canon-ang-pi-3pi)
  then show ?thesis
    by force
qed
have jj:<B ∉ line C A› <Q ∉ line A C› <R ∉ line A B›
  using assms f3' ‹A ≠ C ∧ Q ∉ line A C› ‹A ≠ B ∧ R ∉ line A B›
  by (simp-all add: collinear-sym1 collinear-sym2 line-def)
then have inf-pi3:<abs ?α < pi/3› <abs ?β < pi/3› <abs ?γ < pi/3›
  using ang-inf-pi3[of B C A ?α] ang-inf-pi3[of C A B ?β] ang-inf-pi3[of A B C
?γ]
  by (auto simp add:collinear-sym2 collinear-sym1 assms line-def)
then have <abs (?α + ?β + ?γ) < pi>
  by argo
have j1:<|?α + ?β + ?γ| = pi/3 ⟹ ?a1 * ?a2 * ?a3 = root3›
proof -
  assume hh:<|?α + ?β + ?γ| = pi/3›
  have f20:<cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c C B A / 3)) * cis
(2 * (angle-c A C B / 3)) = cis (2 * (?α + ?β + ?γ))›
    by (simp add: cis-mult)
  then have f21:<cis (2 * (?α + ?β + ?γ)) = cis (2 * (?α + ?β + ?γ))›
    by (smt (verit, ccfv-threshold) canon-ang-cos canon-ang-sin canon-ang-sum
cis.code)
  with hh show ?thesis
    by (metis (no-types, opaque-lifting) f21 f20 times-divide-eq-right root3-def)

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qed
have ⟨rot1.r ∘ rot1.r = complex-rotation.r A (4 * ?α)⟩
  using composed-rotations-same-center by auto
have g10:⟨((rot1.r ∘ rot1.r) ∘ (rot2.r ∘ rot2.r) ∘ (rot3.r ∘ rot3.r)) =
complex-rotation.r A (6 * ?α) ∘ complex-rotation.r B (6 * ?β) ∘ complex-rotation.r C (6 * ?γ)⟩
  using composed-rotations-same-center by auto
have f26:⟨|6 * ?α + 6 * ?β + 6 * ?γ| = 0⟩
  by (smt (verit, ccfv-SIG) canon-ang-sum eq-pi two-pi-canonical)
also have ⟨|6 * ?γ| ≠ 0⟩
proof (rule ccontr)
  assume ⟨¬ |6 * (angle-c A C B / 3)| ≠ 0⟩
  then obtain k::int where ⟨6 * |(angle-c A C B / 3)| = 2 * k * pi⟩
    by (metis add.commute add.inverse-neutral add-0 canon-ang(3)
      diff-conv-add-uminus f10 f13 of-int-mult of-int-numeral)
  then have ⟨3 * |(angle-c A C B / 3)| = k * pi⟩
    by force
  then show False
    by (metis (no-types, opaque-lifting) assms(1) collinear-sin-neq-0 divide-eq-eq-numeral1(1)
      f10 f13 mult-1 sin-kpi times-divide-eq-left zero-neq-numeral)
qed
then have f24:⟨|6 * ?α + 6 * ?β| ≠ 0⟩
  by (metis add.commute add.right-neutral arg-cis calculation
    canon-ang-cos canon-ang-sin canon-ang-sum cis.code)
let ?A1 = ⟨(A * (1 - cis (6 * ?α)) + B * cis (6 * ?α) * (1 - cis (6 * ?β))) / (1 - cis (6 * ?β + 6 * ?α))⟩
have g11:⟨complex-rotation.r A (6 * ?α) ∘ complex-rotation.r B (6 * ?β) ∘ complex-rotation.r C (6 * ?γ) =
complex-rotation.r ?A1 (6 * ?α + 6 * ?β) ∘ complex-rotation.r C (6 * ?γ)⟩
  using composed-rotations[of ⟨6 * ?α⟩ ⟨6 * ?β⟩ A B, OF f24]
  by argo
then have f27:⟨complex-rotation.r ?A1 (6 * ?α + 6 * ?β) ∘ complex-rotation.r C (6 * ?γ) =
(λz. z + (C -
(A * (1 - cis (6 * (angle-c B A C / 3))) + B * cis (6 * (angle-c B A C / 3))) * (1 - cis (6 * (angle-c C B A / 3)))) /
(1 - cis (6 * (angle-c C B A / 3)) + 6 * (angle-c B A C / 3))) * (cis (6 * (angle-c B A C / 3)) + 6 * (angle-c C B A / 3)) - 1))⟩ (is ⟨?l =
(λz. z + ?A2)⟩)
  using composed-rotation-is-trans[of ⟨6 * ?α + 6 * ?β⟩ ⟨6 * ?γ⟩ ?A1 C, OF f26]
  by auto
have f30:⟨2 * angle-c A C B = 6 * ?γ⟩
  by auto
have have f30:⟨2 * angle-c A C B = 6 * ?γ⟩
  have ⟨axial-symmetry C B A = complex-rotation.r C (|2 * angle-c A C B|) A⟩
    using img-r-sym collinear-sym1 f2 jj(1) line-def by auto
  then have first:⟨complex-rotation.r C (6 * ?γ) A = axial-symmetry C B A⟩ (is ⟨?rr = ?axA⟩)
    using canon-ang-cos canon-ang-sin cis.code f30 complex-rotation.r-def by presburger
  have f31:⟨axial-symmetry B C ?axA = complex-rotation.r B (|2 * angle-c ?axA| B

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C]) ?axA
  by (metis ‹A ≠ C ∧ Q ∉ line A C› ‹|6 * (angle-c A C B / 3)| ≠ 0›
    ‹complex-rotation.r C (6 * (angle-c A C B / 3)) A = axial-symmetry C B
A›
    complex-rotation.ang-eq-theta angle-c-neq0
    axial-symmetry-eq f2 img-r-sym involution-symmetry line-is-inv z1-inv)
  have f32:⟨angle-c C B A = 3*?β⟩
    by simp
  then have f33:⟨angle-c ?axA B C = 3*?β⟩
    by (smt (verit) ‹A ≠ B ∧ R ∉ line A B› ‹A ≠ C ∧ Q ∉ line A C› ‹|6 *
      (angle-c A C B / 3)| ≠ 0›
      ‹complex-rotation.r C (6 * (angle-c A C B / 3)) A = axial-symmetry C B
A›
      complex-rotation.ang-eq-theta angle-c-commute angle-c-neq0 angle-symmetry-eq
      assms(1)
      axial-symmetry-eq collinear-sin-neq-0 collinear-sym1 f2 line-is-inv sin-pi)
  then have snd:⟨complex-rotation.r B (6*?β) ((axial-symmetry B C A)) = A›
    by (smt (verit, ccfv-SIG) ‹A ≠ C ∧ Q ∉ line A C› ‹|6 * (angle-c A C B / 3)| ≠ 0›
      ‹complex-rotation.r C (6 * (angle-c A C B / 3)) A = axial-symmetry C B
A›
      complex-rotation.ang-eq-theta angle-c-neq0 axial-symmetry-eq canon-ang-cos
      canon-ang-sin
      cis.code f2 img-r-sym involution-symmetry line-is-inv complex-rotation.r-def
      z1-inv)
  then have thrd:⟨complex-rotation.r A (6*?α) A = A›
    unfolding complex-rotation.r-def by auto
  then have ‹(complex-rotation.r A (6*?α) o complex-rotation.r B (6*?β) o com-
    plex-rotation.r C (6*?γ)) A = A›
    apply(simp)
    by (smt (verit, best) axial-symmetry-eq f30 f32 first snd)
  then have g21:⟨((rot1.r o rot1.r o rot1.r) o (rot2.r o rot2.r o rot2.r) o (rot3.r o rot3.r o rot3.r)) =
  (λz. z + ?A2)›
    apply(subst g10) apply(subst g11) apply(subst f27) by(simp add:fun-eq-iff)
  then have ‹?A2 = 0›
    by (metis ‹(complex-rotation.r A (6 * (angle-c B A C / 3)) o
      complex-rotation.r B (6 * (angle-c C B A / 3)) o
      complex-rotation.r C (6 * (angle-c A C B / 3))) A = A›
      add-diff-cancel-left' cancel-comm-monoid-add-class.diff-cancel g10)
  then have g22:⟨((rot1.r o rot1.r o rot1.r) o (rot2.r o rot2.r o rot2.r) o (rot3.r o rot3.r o rot3.r)) =
  id›
    using g21 by(auto simp:fun-eq-iff)
  have g20:⟨rot1.r z = ?a1*z + ?b1› ⟨rot2.r z = ?a2*z + ?b2› ⟨rot3.r z = ?a3 *
  z + ?b3› for z
    by(auto simp:field-simps complex-rotation.r-def)
  then have ‹((rot1.r o rot1.r o rot1.r) o (rot2.r o rot2.r o rot2.r) o (rot3.r o rot3.r o rot3.r)) =
  (cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c C B A / 3))
  * cis (2 * (angle-c A C B / 3))) ^ 3 * z +
  ((cis (2 * (angle-c B A C / 3)))^2 + cis (2 * (angle-c B A C / 3)) + 1) * (A

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* (1 - cis (2 * (angle-c B A C / 3)))) +
  cis (2 * (angle-c B A C / 3)) ^ 3 * ((cis (2 * (angle-c C B A / 3)))^2 + cis (2
* (angle-c C B A / 3)) + 1) *
  (B * (1 - cis (2 * (angle-c C B A / 3)))) +
  cis (2 * (angle-c B A C / 3)) ^ 3 * cis (2 * (angle-c C B A / 3)) ^ 3 *
  ((cis (2 * (angle-c A C B / 3)))^2 + cis (2 * (angle-c A C B / 3)) + 1) *
  (C * (1 - cis (2 * (angle-c A C B / 3)))) > (is <l z = ?A3 z> for z
  using equality-for-comp[of rot3.r ?a3 ?b3 rot2.r ?a2 ?b2 rot1.r ?a1 ?b1, OF
g20(3) g20(2) g20(1)]
  by blast
  then have <z = ?A3 z> for z
  by (metis g22 id-apply)
  then have very-imp:<((cis (2 * (angle-c B A C / 3)))^2 + cis (2 * (angle-c B A
C / 3)) + 1) * (A * (1 - cis (2 * (angle-c B A C / 3)))) +
  cis (2 * (angle-c B A C / 3)) ^ 3 * ((cis (2 * (angle-c C B A / 3)))^2 + cis (2
* (angle-c C B A / 3)) + 1) *
  (B * (1 - cis (2 * (angle-c C B A / 3)))) +
  cis (2 * (angle-c B A C / 3)) ^ 3 * cis (2 * (angle-c C B A / 3)) ^ 3 *
  ((cis (2 * (angle-c A C B / 3)))^2 + cis (2 * (angle-c A C B / 3)) + 1) *
  (C * (1 - cis (2 * (angle-c A C B / 3)))) = 0>
  by (metis comm-monoid-add-class.add-0 mult-zero-class.mult-zero-right)
{assume hhh:<|?α+?β+?γ|=pi/3>
have j5:<?a1*?a2 ≠ 1>
proof(rule ccontr)
assume <¬ cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c C B A / 3)) ≠
1>
then have <?a1*?a2 = cis 0>
  using cis-zero by presburger
then have <|2*(?α+?β)| = 0>
  by (metis arg-cis cis-mult distrib-left zero-canonical)
then have t:<|?α+?β|=pi ∨ |?α+?β|=0>
  by (smt (verit, del-insts) canon-ang(1) canon-ang-id canon-ang-sum canon-ang-uminus)
then have <|?α+?β|=0 ==> False>
  by (metis add-0 assms(1) canon-ang-sum collinear-sin-neq-0
      divide-cancel-right f10 f13 hhh sin-pi zero-neq-numeral)
then have <|1/3*(3*?α+3*?β)|=pi> using t by(auto simp:algebra-simps)
then have <|(pi-|3*?γ|)| = |(3*?α+3*?β)|>
  by (smt (verit, ccfv-SIG) canon-ang-diff eq-pi)
then have <|(pi/3-|?γ|)| = pi>
  by (smt (verit, best) <|angle-c B A C / 3 + angle-c C B A / 3|=0 ==>
False>
  canon-ang-diff hhh t)
then have <|?γ| ∈ {0<..<pi/3}>
  by (smt (verit, ccfv-SIG) add-divide-distrib ang-vec-bounded angle-c-def
      arccos-minus-one-half arccos-one-half arcsin-minus-one-half arcsin-one-half
      arcsin-plus-arccos canon-ang-id canon-ang-pi-3pi divide-cancel-right f13
      field-sum-of-halves)
  then show False

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using ‹|pi / 3 - |angle-c A C B / 3| = pi› add-divide-distrib canon-ang-id
by auto
qed
have r-eq-all:‹rot1.r z = ?a1*z + ?b1› ‹rot2.r z = ?a2*z + ?b2› ‹rot3.r z =
?a3*z + ?b3› for z
  by(auto simp:field-simps complex-rotation.r-def cis.code)
have f21:‹rot2.r (rot3.r P) = ?a2*?a3*P + ?a2*?b3 + ?b2›
  apply(subst r-eq-all(2)) apply(subst r-eq-all(3)) by(auto simp:field-simps)

then have ‹rot2.r (rot3.r P) = ?a2*?a3*P + ?a2*?b3 + ?b2 ↔ P*(1 - ?a2*?a3)
= ?a2*?b3 + ?b2›
  by (smt (verit) add.commute add-diff-cancel-left' add-diff-eq f15
f16 mult.commute mult.right-neutral right-diff-distrib')
then have j2:‹?a2*?a3 ≠ 1 ⇒ P = (?a2*?b3 + ?b2)/(1 - ?a2*?a3)›
  using f21 by(auto simp:field-simps)
have j4:‹?a1*?a3 ≠ 1›
proof(rule ccontr)
assume ‹¬ cis (2 * (angle-c B A C / 3)) * cis (2 * (angle-c A C B / 3)) ≠
1›
then have ‹?a1*?a3 = cis 0›
  using cis-zero by presburger
then have ‹|2*(?α+?γ)| = 0›
  by (metis arg-cis cis-mult distrib-left zero-canonical)
then have t:‹|?α+?γ| = pi ∨ |?α+?γ| = 0›
  by (smt (verit, del-insts) canon-ang(1) canon-ang-id canon-ang-sum canon-ang-uminus)
then have k0:‹|?α+?γ| = 0 ⇒ False›
  by (smt (verit, ccfv-SIG) ‹A ≠ B ∧ R ∉ line A B› ‹A ≠ C ∧ Q ∉ line A
C›
    ‹|angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3| < pi›
    ‹angle-c (axial-symmetry A B R) A B = angle-c B A C / 3› ‹angle-c R
A B = - (angle-c B A C / 3)›
    ‹angle-c R B A = angle-c C B A / 3› ang-pos-pos assms(3) canon-ang-id
f2 f30 f32 neq-0 z1-inv)

then have ‹|1/3*(3*?α+3*?γ)| = pi› using t by(auto simp:algebra-simps)
then have ‹|(pi - |3*?β|)| = |(3*?α+3*?γ)|›
  by (smt (verit, ccfv-SIG) canon-ang-diff eq-pi)
then have ‹|(pi/3 - |?β|)| = pi›
  by (smt (verit, best) k0
    canon-ang-diff hhh t)
then have k1:‹|?β| ∈ {0 <.. < pi/3}›
  by (smt (verit, del-insts) arccos-one-half arcsin-one-half arcsin-plus-arccos
canon-ang-id divide-nonneg-pos field-sum-of-halves inf-pi3(2) pi-ge-zero)
then show False
  using k1 add-divide-distrib canon-ang-id ‹|pi / 3 - |angle-c C B A / 3| = pi› by auto
qed
have j3:‹?a2*?a3 ≠ 1›
proof(rule ccontr)

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assume  $\neg cis(2 * (angle-c C B A / 3)) * cis(2 * (angle-c A C B / 3)) \neq$ 
1> then have  $\langle ?a2 * ?a3 = cis 0 \rangle$ 
    using cis-zero by presburger
then have  $\langle \lfloor 2 * (?\beta + ?\gamma) \rfloor = 0 \rangle$ 
    by (metis arg-cis cis-mult distrib-left zero-canonical)
then have  $t: \langle \lfloor ?\beta + ?\gamma \rfloor = pi \vee \lfloor ?\beta + ?\gamma \rfloor = 0 \rangle$ 
    by (smt (verit, del-insts) canon-ang(1) canon-ang-id canon-ang-sum canon-ang-uminus)
then have  $k0: \langle \lfloor ?\beta + ?\gamma \rfloor = 0 \implies False \rangle$ 
    by (smt (verit, ccfv-SIG)  $\langle A \neq B \wedge R \notin line A B \rangle \langle A \neq C \wedge Q \notin line A$ 
C>
 $\langle |angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3| < pi \rangle$ 
 $\langle angle-c (axial-symmetry A B R) A B = angle-c B A C / 3 \rangle \langle angle-c R$ 
 $A B = - (angle-c B A C / 3) \rangle$ 
 $\langle angle-c R B A = angle-c C B A / 3 \rangle$  ang-pos-pos assms(3) canon-ang-id
f2 f30 f32 neq-0 z1-inv)

then have  $\langle \lfloor 1/3 * (3 * ?\beta + 3 * ?\gamma) \rfloor = pi \rangle$  using t by (auto simp:algebra-simps)
then have  $\langle \lfloor (pi - \lfloor 3 * ?\alpha \rfloor) \rfloor = \lfloor (3 * ?\beta + 3 * ?\gamma) \rfloor \rangle$ 
    by (smt (verit, ccfv-SIG) canon-ang-diff eq-pi)
then have  $\langle \lfloor (pi/3 - \lfloor ?\alpha \rfloor) \rfloor = pi \rangle$ 
    by (smt (verit, best) k0
        canon-ang-diff hhh t)
then have  $k1: \langle \lfloor ?\alpha \rfloor \in \{0, \dots, pi/3\} \rangle$ 
    by (smt (verit, del-insts) arccos-one-half arcsin-one-half arcsin-plus-arccos
        canon-ang-id divide-nonneg-pos field-sum-of-halves inf-pi3(1) pi-ge-zero)
then show False
    using k1 add-divide-distrib canon-ang-id  $\langle \lfloor (pi/3 - \lfloor ?\alpha \rfloor) \rfloor = pi \rangle$  by auto
qed
have f21':  $\langle rot3.r (rot1.r Q) = ?a3 * ?a1 * Q + ?a3 * ?b1 + ?b3 \rangle$ 
    apply(subst r-eq-all(1)) apply(subst r-eq-all(3)) by(auto simp:field-simps)
have f21'':  $\langle rot1.r (rot2.r R) = ?a1 * ?a2 * R + ?a1 * ?b2 + ?b1 \rangle$ 
    apply(subst r-eq-all(1)) apply(subst r-eq-all(2)) by(auto simp:field-simps)
have jjj:  $\langle ?a1 \neq 0 \wedge ?a2 \neq 0 \wedge ?a3 \neq 0 \rangle$ 
    using cis-neq-zero by blast
then have eq-j:  $\langle ?a1 * ?a2 * ?a3 = root3 \rangle$ 
    using j1 hhh by blast
then have P-def:  $\langle P = (?a2 * ?b3 + ?b2) / (1 - ?a2 * ?a3) \rangle$  (is  $\langle -=?P \rangle$ )
    using j2 j3 by blast
have n1:  $\langle 1 - ?a2 * ?a3 = (?a1 - root3) / ?a1 \rangle$ 
    by (smt (verit, ccfv-threshold) cis-divide cis-times-cis-opposite cis-zero
        diff-divide-distrib eq-j minus-real-def mult-1 times-divide-eq-left)
then have P-last:  $\langle P = ?a1 * (?a2 * ?b3 + ?b2) / (?a1 - root3) \rangle$ 
    using P-def jjj by (simp add:)
then have Q-def:  $\langle Q = (?a3 * ?b1 + ?b3) / (1 - ?a3 * ?a1) \rangle$  (is  $\langle -=?Q \rangle$ )
    using f21' j4 Q-inv by(auto simp:field-simps)
have n2:  $\langle 1 - ?a3 * ?a1 = (?a2 - root3) / ?a2 \rangle$ 
    by (smt (verit, ccfv-threshold) cis-divide cis-times-cis-opposite cis-zero
        diff-divide-distrib eq-j minus-real-def mult-1 times-divide-eq-left)

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then have Q-last: $\langle Q = ?a2*(?a3*b1 + ?b3)/(?a2-root3) \rangle$ 
  using Q-def jjj by simp
then have R-def: $\langle R = (?a1*b2 + ?b1)/(1-?a1*a2) \rangle$  (is  $\leftarrow=?R$ )
  using f21'' j5 R-inv by(auto simp:field-simps)
have n3: $\langle 1-?a1*a2 = (?a3 - root3)/?a3 \rangle$ 
  by (smt (verit, ccfv-threshold) cis-divide cis-times-cis-opposite cis-zero
    diff-divide-distrib eq-j minus-real-def mult-1 times-divide-eq-left)
then have R-last: $\langle R = ?a3*(?a1*b2 + ?b1)/(?a3-root3) \rangle$ 
  using R-def jjj by simp
have  $\langle ?a1 - root3 \neq 0 \wedge ?a2 - root3 \neq 0 \wedge ?a3 - root3 \neq 0 \rangle$ 
  using eq-j j3 j4 j5 by auto
have rule-s: $\langle c \neq 0 \implies a/c + b/c = (a+b)/c \rangle$  for a b c::real
  by argo
define j' where
  defs:  $\langle j' \equiv root3 \rangle$ 
have simp-rules-for-eq: $\langle P = (?a2*b3 + ?b2)/(1-?a2*a3) \rangle$   $\langle R = (?a1*b2 + ?b1)/(1-?a1*a2) \rangle$ 
+  $\langle Q = (?a3*b1 + ?b3)/(1-?a1*a3) \rangle$   $\langle 1-?a1*a2 \neq 0 \wedge 1-?a2*a3 \neq 0 \wedge 1-?a1*a3 \neq 0 \rangle$ 
 $\langle ?a1*a2*a3 = j' \rangle$ 
 $\langle 1 + j' + j'^2 = 0 \rangle$   $\langle ((1 - ?a1 * ?a3) * ((1 - ?a1 * ?a2) * (1 - ?a2 * ?a3))) \neq 0 \rangle$ 
 $\langle j'^3 = 1 \rangle$   $\langle j'*j'*j' = 1 \rangle$   $\langle ?a1 \neq 0 \rangle$   $\langle ?a2 \neq 0 \rangle$   $\langle ?a3 \neq 0 \rangle$ 
 $\langle (1 - ?a1 * ?a2)*?a3 = (?a3 - j') \rangle$   $\langle (1 - ?a2 * ?a3)*?a1 = (?a1 - j') \rangle$   $\langle (1 - ?a1 * ?a3)*?a2 = (?a2 - j') \rangle$ 
= using Q-def P-def R-def eq-j j3 jjj j4 j5 root3-eq-0 root-unity-3 n1 n2 n3
  by(auto simp add:mult.commute power3-eq-cube defs root3-def)
have graal: $\langle (?a1^2 + ?a1 + 1)*?b1 + ?a1^3*(?a2^2 + ?a2 + 1)*?b2 + ?a1^3*?a2^3*(?a3^2 + ?a3 + 1)*?b3 = -j'*?a1^2*?a2*(?a1 - j')*(?a2 - j')*(?a3 - j')*(?R + j'*?P + j'^2*?Q) \rangle$ 
  using root-unity-carac[of ?a1 ?a2 ?a3 j' ?R ?b2 ?b1 ?P ?b3 ?Q]
  using simp-rules-for-eq defs by(auto simp:)
then have  $\langle (?a1^2 + ?a1 + 1)*?b1 + ?a1^3*(?a2^2 + ?a2 + 1)*?b2 + ?a1^3*?a2^3*(?a3^2 + ?a3 + 1)*?b3 = 0 \rangle$ 
  unfolding defs using very-imp by auto
then have  $\langle (?R + j'*?P + j'^2*?Q) = 0 \rangle$ 
  using graal simp-rules-for-eq(13) simp-rules-for-eq(14)
  simp-rules-for-eq(15) simp-rules-for-eq(11) simp-rules-for-eq(12) simp-rules-for-eq(4)
by force
then have impp: $\langle ?R + j'*?P + j'^2*?Q = 0 \rangle$ 
  using R-def P-def Q-def
  by (metis simp-rules-for-eq(1) simp-rules-for-eq(2))
have neq-all: $\langle A \neq B \wedge B \neq C \wedge C \neq A \rangle$ 
  using  $\langle A \neq B \wedge R \notin \text{line } A B \rangle$   $\langle A \neq C \wedge Q \notin \text{line } A C \rangle$  f2 by argo
have  $\langle R \neq Q \rangle$ 
proof (rule ccontr)
  assume  $\neg R \neq Q$ 
then have q1: $\langle \text{angle-c } A B R = \text{angle-c } A B Q \rangle$ 
   $\langle \text{angle-c } B A R = \text{angle-c } B A Q \rangle$   $\langle \text{angle-c } C A Q = \text{angle-c } C A R \rangle$ 

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⟨angle-c A C Q = angle-c A C R⟩
  using assms by auto
then have ⟨angle-c B A R = ?α⟩ using assms by auto
then have ⟨angle-c B A Q = ?α⟩
  using q1 by auto
then have ⟨angle-c A B Q = angle-c A B R + angle-c R B Q⟩
  using ⟨¬ R ≠ Q⟩ angle-c-neq0 by auto
have ⟨angle-c C A B = - 3 * ?α⟩ using assms
  using ⟨angle-c C A Q = - (angle-c B A C / 3)⟩ by argo
then have ⟨angle-c (axial-symmetry B A R) A B = ?α⟩
  by (metis ⟨angle-c (axial-symmetry A B R) A B = angle-c B A C / 3⟩
  axial-symmetry-eq)
have ⟨C - A ≠ 0 ∧ Q - A ≠ 0 ∧ A - R ≠ 0 ∧ A - B ≠ 0⟩
  using ⟨A ≠ C ∧ Q ≠ line A C⟩ ⟨angle-c A B Q = angle-c A B R + angle-c
R B Q⟩
  ⟨angle-c R B A = angle-c C B A / 3⟩ neq-0 q1(1) neq-all by fastforce
  then have ⟨angle-c C A B = angle-c C A Q + angle-c Q A R + angle-c R
A B⟩
    using angle-c-sum[of C A Q R] angle-c-sum[of C A R B]
    by (smt (verit) ⟨¬ R ≠ Q⟩ ⟨angle-c C A B = - 3 * (angle-c B A C / 3)⟩
    ⟨angle-c C A Q = - (angle-c B A C / 3)⟩ ⟨angle-c R A B = - (angle-c
B A C / 3)⟩
      angle-c-neq0 canon-ang(1) canon-ang(2) canon-ang-id)
  then show False
    using ⟨¬ R ≠ Q⟩
    by (smt (verit, best) ⟨angle-c C A B = - 3 * (angle-c B A C / 3)⟩
    ⟨angle-c C A Q = - (angle-c B A C / 3)⟩ ⟨angle-c R A B = - (angle-c
B A C / 3)⟩
      angle-c-neq0 neq-0 zero-canonical)
qed
then have ⟨cdist R P = cdist R Q ∧ cdist R Q = cdist Q P⟩
  using equilateral-characterization-neg[of R Q P] impp unfolding defs
  by (metis cdist-commute)
then have ?thesis
  using cdist-commute by auto
}note case-pi3=this
then show ?thesis using hhh by auto
qed

```

theorem Morley-neg:

assumes ¬collinear A B C

⟨angle-c A B R = angle-c A B C / 3⟩ (is ⟨?abr = ?abc⟩)

⟨angle-c B A R = angle-c B A C / 3⟩ (is ⟨?bar = ?α⟩)

⟨angle-c B C P = angle-c B C A / 3⟩ (is ⟨?bcp = ?bca⟩)

⟨angle-c C B P = angle-c C B A / 3⟩ (is ⟨?cbp = ?β⟩)

⟨angle-c C A Q = angle-c C A B / 3⟩ (is ⟨?caq = ?cab⟩)

⟨angle-c A C Q = angle-c A C B / 3⟩ (is ⟨?acq = ?γ⟩)

and hhh: ⟨angle-c B A C / 3 + angle-c C B A / 3 + angle-c A C B / 3⟩ =
-pi/3

shows $\langle cdist R P = cdist P Q \wedge cdist Q R = cdist P Q \rangle$
proof –
have $\langle |-angle-c B A C / 3 + -angle-c C B A / 3 + -angle-c A C B / 3 \rangle = pi/3 \rangle$
using hhh
by (metis (no-types, opaque-lifting) canon-ang(1) canon-ang-uminus less-eq-real-def

*linorder-not-le minus-add-distrib minus-divide-left mult-le-0-iff
 nonzero-mult-div-cancel-left not-numeral-le-zero pi-gt-zero times-divide-eq-right
 verit-minus-simplify(4) zero-canonical*

then show ?thesis **using** Morley-pos[of $C B A P Q R$]
by (smt (verit, best) Morley-pos angle-c-commute assms(1) assms(2) assms(3)
assms(4) assms(5)
assms(6) assms(7) cdist-commute collinear-sin-neq-0 collinear-sym1 collinear-sym2
sin-pi)
qed

theorem Morley:
assumes $\neg collinear A B C$
 $\langle angle-c A B R = angle-c A B C / 3 \rangle$ (**is** $\langle ?abr = ?abc \rangle$)
 $\langle angle-c B A R = angle-c B A C / 3 \rangle$ (**is** $\langle ?bar = ?\alpha \rangle$)
 $\langle angle-c B C P = angle-c B C A / 3 \rangle$ (**is** $\langle ?bcp = ?bca \rangle$)
 $\langle angle-c C B P = angle-c C B A / 3 \rangle$ (**is** $\langle ?cbp = ?\beta \rangle$)
 $\langle angle-c C A Q = angle-c C A B / 3 \rangle$ (**is** $\langle ?caq = ?cab \rangle$)
 $\langle angle-c A C Q = angle-c A C B / 3 \rangle$ (**is** $\langle ?acq = ?\gamma \rangle$)
shows $\langle cdist R P = cdist P Q \wedge cdist Q R = cdist P Q \rangle$
proof –
{fix $A B C \gamma$
assume ABC-nline: $\langle A \notin line B C \rangle$
and eq-3c: $\langle angle-c A C B = 3*\gamma \rangle$
then have neq-PI: $\langle abs \gamma < pi/3 \rangle$
proof –
have $\langle angle-c A C B \neq pi \rangle$ **using** ABC-nline(1) unfolding line-def
by (metis angle-c-commute-pi collinear-iff mem-Collect-eq non-collinear-independant)
then have $\langle angle-c A C B \in \{-pi..pi\} \rangle$
using ang-c-in less-eq-real-def **by** auto
then show $\langle abs \gamma < pi/3 \rangle$
using eq-3c **by** force
qed}
note ang-inf-pi3=this

have $\langle |angle-c B A C + angle-c C B A + angle-c A C B| = pi \rangle$
by (metis Elementary-Complex-Geometry.collinear-def add.commute
angle-sum-triangle-c assms(1) collinear-sym1 collinear-sym2')
have eq-pi: $\langle 3*\alpha + 3*\beta + 3*\gamma \rangle = pi \rangle$
using $\langle |angle-c B A C + angle-c C B A + angle-c A C B| = pi \rangle$ **by** force
have $\langle A \notin line B C \wedge B \notin line A C \wedge C \notin line A B \rangle$
using assms(1) unfolding line-def **using** collinear-sym1 collinear-sym2' **by**(auto)

then have f2: $\langle abs \alpha < pi/3 \wedge abs \beta < pi/3 \wedge abs \gamma < pi/3 \rangle$
using ang-inf-pi3
by (smt (verit, ccfv-SIG) ang-vec-bounded angle-c-def assms(1) collinear-sin-neq-0

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collinear-sym1 collinear-sym2 divide-less-cancel minus-divide-left sin-pi)
have possi-abc:  $\lfloor \alpha + \beta + \gamma \rfloor = \pi/3 \vee \lfloor \alpha + \beta + \gamma \rfloor = -\pi/3$ 
proof -
  have  $\alpha + \beta + \gamma \leq \pi$ 
  using f2 by(auto simp:field-simps)
  then have  $\lfloor \alpha + \beta + \gamma \rfloor = \alpha + \beta + \gamma$ 
    using canon-ang-id f2 by fastforce
  then have  $\lfloor 3 * \lfloor \alpha + \beta + \gamma \rfloor \rfloor = \pi$ 
    using eq-pi by force
  then show ?thesis
    by (smt (verit) add-divide-distrib arccos-minus-one-half arccos-one-half arc-
sin-minus-one-half
      arcsin-one-half arcsin-plus-arccos canon-ang-id canon-ang-minus-3pi-minus-pi canon-ang-pi-3pi
      eq-pi f2 field-sum-of-halves)
  qed
  then show ?thesis
    using Morley-neg Morley-pos assms(1) assms(2) assms(3) assms(4) assms(5)
    assms(6) assms(7)
    by meson
  qed
end

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References

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- [3] Wiedijk's catalogue "Formalizing 100 Theorems" <https://www.cs.ru.nl/~freek/100/>, It appears at position 121....