

# On the Formalization of Martingales

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## Abstract

In the scope of this project, we present a formalization of martingales in arbitrary Banach spaces using Isabelle/HOL.

The current formalization of conditional expectation in the Isabelle library is limited to real-valued functions. To overcome this limitation, we extend the construction of conditional expectation to general Banach spaces, employing an approach similar to the one described in [1]. We use measure theoretic arguments to construct the conditional expectation using suitable limits of simple functions.

Subsequently, we define stochastic processes and introduce the concepts of adapted, progressively measurable and predictable processes using suitable locale definitions<sup>1</sup>. We show the relation

$$\text{adapted} \supseteq \text{progressive} \supseteq \text{predictable}$$

Furthermore, we show that progressive measurability and adaptedness are equivalent when the indexing set is discrete. We pay special attention to predictable processes in discrete-time, showing that  $(X_n)_{n \in \mathbb{N}}$  is predictable if and only if  $(X_{n+1})_{n \in \mathbb{N}}$  is adapted.

Moving forward, we rigorously define martingales, submartingales, and supermartingales, presenting their first consequences and corollaries<sup>2</sup>. Discrete-time martingales are given special attention in the formalization. In every step of our formalization, we make extensive use of the powerful locale system of Isabelle.

The formalization further contributes by generalizing concepts in Bochner integration by extending their application from the real numbers to arbitrary Banach spaces equipped with a second-countable topology. Induction schemes for integrable simple functions on Banach spaces are introduced, accommodating various scenarios with or without a real vector ordering<sup>3</sup>. Specifically, we formalize a powerful result called the “Averaging Theorem”[3] which allows us to show that densities are unique in Banach spaces.

In-depth information on the formalization and the proofs of the individual theorems can be found in [2].

---

<sup>1</sup>[Martingale.Stochastic\\_Process](#)

<sup>2</sup>[Martingale.Martingale](#)

<sup>3</sup>[Martingale.Bochner\\_Integration\\_Addendum](#)

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```

theory Measure-Space-Supplement
imports HOL-Analysis.Measure-Space
begin

```

## 1 Supplementary Lemmas for Measure Spaces

### 1.1 $\sigma$ -Algebra Generated by a Family of Functions

**definition** family-vimage-algebra :: ' $a$  set  $\Rightarrow$  (' $a \Rightarrow$  ' $b$ ) set  $\Rightarrow$  ' $b$  measure  $\Rightarrow$  ' $a$  measure **where**  

$$\text{family-vimage-algebra } \Omega S M \equiv \text{sigma } \Omega (\bigcup_{f \in S. \{f -` A \cap \Omega | A. A \in M\}})$$

For singleton  $S$ , i.e.  $S = \{f\}$  for some  $f$ , the definition simplifies to that of *vimage-algebra*.

**lemma** family-vimage-algebra-singleton: family-vimage-algebra  $\Omega \{f\} M = \text{vimage-algebra } \Omega f M$  **unfolding** family-vimage-algebra-def vimage-algebra-def **by** simp

**lemma**  
**shows** sets-family-vimage-algebra: sets (family-vimage-algebra  $\Omega S M) = \text{sigma-sets } \Omega (\bigcup_{f \in S. \{f -` A \cap \Omega | A. A \in M\}}$   
**and** space-family-vimage-algebra[simp]: space (family-vimage-algebra  $\Omega S M) = \Omega$   
**by** (auto simp add: family-vimage-algebra-def sets-measure-of-conv space-measure-of-conv)

**lemma** measurable-family-vimage-algebra:  
**assumes**  $f \in S f \in \Omega \rightarrow \text{space } M$   
**shows**  $f \in \text{family-vimage-algebra } \Omega S M \rightarrow_M M$   
**using** assms **by** (intro measurableI, auto simp add: sets-family-vimage-algebra)

**lemma** measurable-family-vimage-algebra-singleton:  
**assumes**  $f \in \Omega \rightarrow \text{space } M$   
**shows**  $f \in \text{family-vimage-algebra } \Omega \{f\} M \rightarrow_M M$   
**using** assms measurable-family-vimage-algebra **by** blast

A collection of functions are measurable with respect to some  $\sigma$ -algebra  $N$ , if and only if the  $\sigma$ -algebra they generate is contained in  $N$ .

**lemma** measurable-family-iff-sets:  
**shows**  $(S \subseteq N \rightarrow_M M) \longleftrightarrow S \subseteq \text{space } N \rightarrow \text{space } M \wedge \text{family-vimage-algebra } (\text{space } N) S M \subseteq N$   
**proof** (standard, goal-cases)  
**case** 1  
**hence** subset:  $S \subseteq \text{space } N \rightarrow \text{space } M$  **using** measurable-space **by** fast  
**have**  $\{f -` A \cap \text{space } N | A. A \in M\} \subseteq N$  **if**  $f \in S$  **for**  $f$  **using** measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric], off] 1 subset that  
**by** (fastforce simp add: sets-family-vimage-algebra)  
**then show** ?case **unfolding** sets-family-vimage-algebra **using** sets.sigma-algebra-axioms  
**by** (simp add: subset, intro sigma-algebra.sigma-sets-subset, blast+)

```

next
  case 2
    hence subset:  $S \subseteq \text{space } N \rightarrow \text{space } M$  by simp
    show ?case
    proof (standard, goal-cases)
      case (1 x)
        have family-vimage-algebra (space N) {x} M ⊆ N by (metis (no-types, lifting)
1 2 sets-family-vimage-algebra SUP-le-iff sigma-sets-le-sets-iff singletonD)
        thus ?case using measurable-iff-sets[unfolded family-vimage-algebra-singleton[symmetric]]
subset[THEN subsetD, OF 1] by fast
      qed
    qed

lemma family-vimage-algebra-diff:
  shows family-vimage-algebra  $\Omega$  S M = sigma  $\Omega$  (sets (family-vimage-algebra  $\Omega$ 
( $S - I$ ) M)  $\cup$  family-vimage-algebra  $\Omega$  ( $S \cap I$ ) M)
  using sets.space-closed space-measure-of-conv
  unfolding family-vimage-algebra-def sets-family-vimage-algebra
  by (intro sigma-eqI, blast, fastforce)
    (intro sigma-sets-eqI, blast, simp add: sets-measure-of-conv split: if-splits,
     meson Diff-subset Sup-subset-mono in-mono inf-sup-ord(1) sigma-sets-subseteq
     subset-image-iff, fastforce+)

end

theory Conditional-Expectation-Banach
imports HOL-Probability.Conditional-Expectation HOL-Probability.Independent-Family

begin

```

## 2 Conditional Expectation in Banach Spaces

While constructing the conditional expectation operator, we have come up with the following approach, which is based on the construction in [1]. Both our approach, and the one in [1] are based on showing that the conditional expectation is a contraction on some dense subspace of the space of functions  $L^1(E)$ . In our approach, we start by constructing the conditional expectation explicitly for simple functions. Then we show that the conditional expectation is a contraction on simple functions, i.e.  $\|E(s|F)(x)\| \leq E(\|s(x)\||F)$  for  $\mu$ -almost all  $x \in \Omega$  with  $s : \Omega \rightarrow E$  simple and integrable. Using this, we can show that the conditional expectation of a convergent sequence of simple functions is again convergent. Finally, we show that this limit exhibits the properties of a conditional expectation. This approach has the benefit of being straightforward and easy to implement, since we could make use of the existing formalization for real-valued functions. To use the construction in [1] we need more tools from functional analysis, which

Isabelle/HOL currently does not have.

Before we can talk about 'the' conditional expectation, we must define what it means for a function to have a conditional expectation.

```
definition has-cond-exp :: 'a measure  $\Rightarrow$  'a measure  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\Rightarrow$  'b::{real-normed-vector, second-countable-topology})  $\Rightarrow$  bool where
  has-cond-exp M F f g = (( $\forall A \in \text{sets } F$ . ( $\int x \in A. f x \partial M$ ) = ( $\int x \in A. g x \partial M$ )
     $\wedge$  integrable M f
     $\wedge$  integrable M g
     $\wedge$  g  $\in$  borel-measurable F)
```

This predicate precisely characterizes what it means for a function  $f$  to have a conditional expectation  $g$ , with respect to the measure  $M$  and the sub- $\sigma$ -algebra  $F$ .

```
lemma has-cond-expI':
  assumes  $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$ 
    integrable M f
    integrable M g
    g  $\in$  borel-measurable F
  shows has-cond-exp M F f g
  using assms unfolding has-cond-exp-def by simp
```

```
lemma has-cond-expD:
  assumes has-cond-exp M F f g
  shows  $\bigwedge A. A \in \text{sets } F \implies (\int x \in A. f x \partial M) = (\int x \in A. g x \partial M)$ 
    integrable M f
    integrable M g
    g  $\in$  borel-measurable F
  using assms unfolding has-cond-exp-def by simp+
```

Now we can use Hilbert's  $\epsilon$ -operator to define the conditional expectation, if it exists.

```
definition cond-exp :: 'a measure  $\Rightarrow$  'a measure  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\Rightarrow$  'b::{banach, second-countable-topology}) where
  cond-exp M F f = (if  $\exists g. \text{has-cond-exp } M F f g$  then (SOME g. has-cond-exp M F f g) else ( $\lambda\_. 0$ ))
```

```
lemma borel-measurable-cond-exp[measurable]: cond-exp M F f  $\in$  borel-measurable F
  by (metis cond-exp-def someI has-cond-exp-def borel-measurable-const)
```

```
lemma integrable-cond-exp[intro]: integrable M (cond-exp M F f)
  by (metis cond-exp-def has-cond-expD(3) integrable-zero someI)
```

```
lemma set-integrable-cond-exp[intro]:
  assumes A  $\in$  sets M
  shows set-integrable M A (cond-exp M F f) using integrable-mult-indicator[OF
```

```
assms integrable-cond-exp, of F f] by (auto simp add: set-integrable-def intro!: integrable-mult-indicator[OF assms integrable-cond-exp])
```

```
lemma has-cond-exp-self:
```

```
assumes integrable M f
shows has-cond-exp M (vimage-algebra (space M) f borel) f f
using assms by (auto intro!: has-cond-expI' measurable-vimage-algebra1)
```

```
lemma has-cond-exp-sets-cong:
```

```
assumes sets F = sets G
shows has-cond-exp M F = has-cond-exp M G
using assms unfolding has-cond-exp-def by force
```

```
lemma cond-exp-sets-cong:
```

```
assumes sets F = sets G
shows AE x in M. cond-exp M F f x = cond-exp M G f x
by (intro AE-I2, simp add: cond-exp-def has-cond-exp-sets-cong[OF assms, of M])
```

```
context sigma-finite-subalgebra
```

```
begin
```

```
lemma borel-measurable-cond-exp'[measurable]: cond-exp M F f ∈ borel-measurable M
by (metis cond-exp-def someI has-cond-exp-def borel-measurable-const subalg measurable-from-subalg)
```

```
lemma cond-exp-null:
```

```
assumes ∄ g. has-cond-exp M F f g
shows cond-exp M F f = (λ-. 0)
unfolding cond-exp-def using assms by argo
```

We state the tower property of the conditional expectation in terms of the predicate *has-cond-exp*.

```
lemma has-cond-exp-nested-subalg:
```

```
fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
assumes subalgebra G F has-cond-exp M F f h has-cond-exp M G f h'
shows has-cond-exp M F h' h
by (intro has-cond-expI') (metis assms has-cond-expD in-mono subalgebra-def)+
```

The following lemma shows that the conditional expectation is unique as an element of L1, given that it exists.

```
lemma has-cond-exp-charact:
```

```
fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
assumes has-cond-exp M F f g
shows has-cond-exp M F f (cond-exp M F f)
AE x in M. cond-exp M F f x = g x
```

```
proof –
```

```

show cond-exp: has-cond-exp M F f (cond-exp M F f) using assms someI
cond-exp-def by metis
let ?MF = restr-to-subalg M F
interpret sigma-finite-measure ?MF by (rule sigma-fin-subalg)
{
  fix A assume A ∈ sets ?MF
  then have [measurable]: A ∈ sets F using sets-restr-to-subalg[OF subalg] by
  simp
  have (ʃ x ∈ A. g x ∂?MF) = (ʃ x ∈ A. g x ∂M) using assms subalg by (auto
  simp add: integral-subalgebra2 set-lebesgue-integral-def dest!: has-cond-expD)
  also have ... = (ʃ x ∈ A. cond-exp M F f x ∂M) using assms cond-exp by
  (simp add: has-cond-exp-def)
  also have ... = (ʃ x ∈ A. cond-exp M F f x ∂?MF) using subalg by (auto simp
  add: integral-subalgebra2 set-lebesgue-integral-def)
  finally have (ʃ x ∈ A. g x ∂?MF) = (ʃ x ∈ A. cond-exp M F f x ∂?MF) by
  simp
}
hence AE x in ?MF. cond-exp M F f x = g x using cond-exp assms subalg by
(intro density-unique-banach, auto dest: has-cond-expD intro!: integrable-in-subalg)
then show AE x in M. cond-exp M F f x = g x using AE-restr-to-subalg[OF
subalg] by simp
qed

```

**corollary** cond-exp-charact:

```

fixes f :: 'a ⇒ 'b:{second-countable-topology, banach}
assumes ⋀A. A ∈ sets F ⇒ (ʃ x ∈ A. f x ∂M) = (ʃ x ∈ A. g x ∂M)
  integrable M f
  integrable M g
  g ∈ borel-measurable F
shows AE x in M. cond-exp M F f x = g x
by (intro has-cond-exp-charact has-cond-expI' assms) auto

```

Identity on F-measurable functions:

If an integrable function  $f$  is already  $F$ -measurable, then  $\text{cond-exp } M F f = f$   $\mu$ -a.e. This is a corollary of the lemma on the characterization of  $\text{cond-exp}$ .

**corollary** cond-exp-F-meas[intro, simp]:

```

fixes f :: 'a ⇒ 'b:{second-countable-topology, banach}
assumes integrable M f
  f ∈ borel-measurable F
shows AE x in M. cond-exp M F f x = f x
by (rule cond-exp-charact, auto intro: assms)

```

Congruence

**lemma** has-cond-exp-cong:

```

assumes integrable M f ⋀x. x ∈ space M ⇒ f x = g x has-cond-exp M F g h
shows has-cond-exp M F f h
proof (intro has-cond-expI'[OF - assms(1)])
  fix A assume asm: A ∈ sets F

```

```

hence set-lebesgue-integral M A f = set-lebesgue-integral M A g by (intro set-lebesgue-integral-cong)
(meson assms(2) subalg in-mono subalgebra-def sets.sets-into-space subalgebra-def
subsetD)+
thus set-lebesgue-integral M A f = set-lebesgue-integral M A h using asm assms(3)
by (simp add: has-cond-exp-def)
qed (auto simp add: has-cond-expD[OF assms(3)])
```

**lemma** cond-exp-cong:

```

fixes f :: 'a ⇒ 'b:{second-countable-topology,banach}
assumes integrable M f integrable M g ∧ x. x ∈ space M ⇒ f x = g x
shows AE x in M. cond-exp M F f x = cond-exp M F g x
proof (cases ∃ h. has-cond-exp M F f h)
  case True
    then obtain h where h: has-cond-exp M F f h has-cond-exp M F g h using
has-cond-exp-cong assms by metis
    show ?thesis using h[THEN has-cond-exp-charact(2)] by fastforce
  next
    case False
    moreover have ∉ h. has-cond-exp M F g h using False has-cond-exp-cong assms
by auto
    ultimately show ?thesis unfolding cond-exp-def by auto
  qed
```

**lemma** has-cond-exp-cong-AE:

```

assumes integrable M f AE x in M. f x = g x has-cond-exp M F g h
shows has-cond-exp M F f h
using assms(1,2) subalg subalgebra-def subset-iff
by (intro has-cond-expI', subst set-lebesgue-integral-cong-AE[OF - assms(1)[THEN
borel-measurable-integrable] borel-measurable-integrable(1)[OF has-cond-expD(2)[OF
assms(3)]]])
(fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF -- AE-symmetric])+
```

**lemma** has-cond-exp-cong-AE':

```

assumes h ∈ borel-measurable F AE x in M. h x = h' x has-cond-exp M F f h'
shows has-cond-exp M F f h
using assms(1, 2) subalg subalgebra-def subset-iff
using AE-restr-to-subalg2[OF subalg assms(2)] measurable-from-subalg
by (intro has-cond-expI', subst set-lebesgue-integral-cong-AE[OF - measurable-from-subalg(1,1)[OF
subalg], OF - assms(1) has-cond-expD(4)[OF assms(3)]])
(fast intro: has-cond-expD[OF assms(3)] integrable-cong-AE-imp[OF -- AE-symmetric])+
```

**lemma** cond-exp-cong-AE:

```

fixes f :: 'a ⇒ 'b:{second-countable-topology,banach}
assumes integrable M f integrable M g AE x in M. f x = g x
shows AE x in M. cond-exp M F f x = cond-exp M F g x
proof (cases ∃ h. has-cond-exp M F f h)
  case True
    then obtain h where h: has-cond-exp M F f h has-cond-exp M F g h using
has-cond-exp-cong-AE assms by (metis (mono-tags, lifting) eventually-mono)
```

```

show ?thesis using h[THEN has-cond-exp-charact(2)] by fastforce
next
  case False
    moreover have  $\nexists h. \text{has-cond-exp } M F g h$  using False has-cond-exp-cong-AE
    assms by auto
    ultimately show ?thesis unfolding cond-exp-def by auto
  qed

```

The conditional expectation operator on the reals, *real-cond-exp*, satisfies the conditions of the conditional expectation as we have defined it.

```

lemma has-cond-exp-real:
  fixes f :: 'a  $\Rightarrow$  real
  assumes integrable M f
  shows has-cond-exp M F f (real-cond-exp M F f)
  by (intro has-cond-expI', auto intro!: real-cond-exp-intA assms)

```

```

lemma cond-exp-real[intro]:
  fixes f :: 'a  $\Rightarrow$  real
  assumes integrable M f
  shows AE x in M. cond-exp M F f x = real-cond-exp M F f x
  using has-cond-exp-charact has-cond-exp-real assms by blast

```

```

lemma cond-exp-cmult:
  fixes f :: 'a  $\Rightarrow$  real
  assumes integrable M f
  shows AE x in M. cond-exp M F ( $\lambda x. c * f x$ ) x = c * cond-exp M F f x
  using real-cond-exp-cmult[OF assms(1), of c] assms(1)[THEN cond-exp-real]
  assms(1)[THEN integrable-mult-right, THEN cond-exp-real, of c] by fastforce

```

## 2.1 Existence

Showing the existence is a bit involved. Specifically, what we aim to show is that *has-cond-exp* M F f (*cond-exp* M F f) holds for any Bochner-integrable f. We will employ the standard machinery of measure theory. First, we will prove existence for indicator functions. Then we will extend our proof by linearity to simple functions. Finally we use a limiting argument to show that the conditional expectation exists for all Bochner-integrable functions.

Indicator functions

```

lemma has-cond-exp-indicator:
  assumes A  $\in$  sets M emeasure M A  $< \infty$ 
  shows has-cond-exp M F ( $\lambda x. \text{indicat-real } A x *_R y$ ) ( $\lambda x. \text{real-cond-exp } M F$ 
  (indicator A) x *_R y)
  proof (intro has-cond-expI', goal-cases)
    case (1 B)
    have  $(\int x \in B. (\text{indicat-real } A x *_R y) \partial M) = (\int x \in B. \text{indicat-real } A x \partial M) *_R$ 
    y using assms by (intro set-integral-scaleR-left, meson 1 in-mono subalg subalgebra-def, blast)

```

```

also have ... = ( $\int x \in B. \text{real-cond-exp } M F (\text{indicator } A) x \partial M$ ) *R y using 1
assms by (subst real-cond-exp-intA, auto)
also have ... = ( $\int x \in B. (\text{real-cond-exp } M F (\text{indicator } A) x *_R y) \partial M$ ) us-
ing assms by (intro set-integral-scaleR-left[symmetric], meson 1 in-mono subalg
subalgebra-def, blast)
finally show ?case .
next
case 2
show ?case using integrable-scaleR-left integrable-real-indicator assms by blast
next
case 3
show ?case using assms by (intro integrable-scaleR-left, intro real-cond-exp-int,
blast+)
next
case 4
show ?case by (intro borel-measurable-scaleR, intro Conditional-Expectation.borel-measurable-cond-exp,
simp)
qed

```

```

lemma cond-exp-indicator[intro]:
fixes y :: 'b::{second-countable-topology,banach}
assumes [measurable]: A ∈ sets M emeasure M A < ∞
shows AE x in M. cond-exp M F (λx. indicat-real A x *R y) x = cond-exp M F
(indicator A) x *R y
proof -
have AE x in M. cond-exp M F (λx. indicat-real A x *R y) x = real-cond-exp M F
(indicator A) x *R y using has-cond-exp-indicator[OF assms] has-cond-exp-charact
by blast
thus ?thesis using cond-exp-real[OF integrable-real-indicator, OF assms] by fast-
force
qed

```

Addition

```

lemma has-cond-exp-add:
fixes f g :: 'a ⇒ 'b::{second-countable-topology,banach}
assumes has-cond-exp M F f f' has-cond-exp M F g g'
shows has-cond-exp M F (λx. f x + g x) (λx. f' x + g' x)
proof (intro has-cond-expI', goal-cases)
case (1 A)
have ( $\int x \in A. (f x + g x) \partial M$ ) = ( $\int x \in A. f x \partial M$ ) + ( $\int x \in A. g x \partial M$ ) using
assms[THEN has-cond-expD(2)] subalg 1 by (intro set-integral-add(2), auto simp
add: subalgebra-def set-integrable-def intro: integrable-mult-indicator)
also have ... = ( $\int x \in A. f' x \partial M$ ) + ( $\int x \in A. g' x \partial M$ ) using assms[THEN
has-cond-expD(1)[OF - 1]] by argo
also have ... = ( $\int x \in A. (f' x + g' x) \partial M$ ) using assms[THEN has-cond-expD(3)]
subalg 1 by (intro set-integral-add(2)[symmetric], auto simp add: subalgebra-def
set-integrable-def intro: integrable-mult-indicator)
finally show ?case .
next

```

```

case 2
show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(2))
next
case 3
show ?case by (metis Bochner-Integration.integrable-add assms has-cond-expD(3))
next
case 4
show ?case using assms borel-measurable-add has-cond-expD(4) by blast
qed

lemma has-cond-exp-scaleR-right:
fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
assumes has-cond-exp M F f f'
shows has-cond-exp M F (λx. c *R f x) (λx. c *R f' x)
using has-cond-expD[OF assms] by (intro has-cond-expI', auto)

lemma cond-exp-scaleR-right:
fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
assumes integrable M f
shows AE x in M. cond-exp M F (λx. c *R f x) x = c *R cond-exp M F f x
proof (cases ?f'. has-cond-exp M F f f')
case True
then show ?thesis using assms has-cond-exp-charact has-cond-exp-scaleR-right
by metis
next
case False
show ?thesis
proof (cases c = 0)
case True
then show ?thesis by simp
next
case c-nonzero: False
have ?f'. has-cond-exp M F (λx. c *R f x) f'
proof (standard, goal-cases)
case 1
then obtain f' where f': has-cond-exp M F (λx. c *R f x) f' by blast
have has-cond-exp M F f (λx. inverse c *R f' x) using has-cond-expD[OF f]
divideR-right[OF c-nonzero] assms by (intro has-cond-expI', auto)
then show ?case using False by blast
qed
then show ?thesis using cond-exp-null[OF False] cond-exp-null by force
qed
qed

lemma cond-exp-uminus:
fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
assumes integrable M f
shows AE x in M. cond-exp M F (λx. - f x) x = - cond-exp M F f x
using cond-exp-scaleR-right[OF assms, of -1] by force

```

Together with the induction scheme *integrable-simple-function-induct*, we can show that the conditional expectation of an integrable simple function exists.

```
corollary has-cond-exp-simple:
  fixes f :: 'a ⇒ 'b::{second-countable-topology, banach}
  assumes simple-function M f emeasure M {y ∈ space M. f y ≠ 0} ≠ ∞
  shows has-cond-exp M F f (cond-exp M F f)
  using assms
proof (induction rule: integrable-simple-function-induct)
  case (cong f g)
    then show ?case using has-cond-exp-cong by (metis (no-types, opaque-lifting)
Bochner-Integration.integrable-cong has-cond-expD(2) has-cond-exp-charact(1))
next
  case (indicator A y)
    then show ?case using has-cond-exp-charact[OF has-cond-exp-indicator] by fast
next
  case (add u v)
    then show ?case using has-cond-exp-add has-cond-exp-charact(1) by blast
qed
```

Now comes the most difficult part. Given a convergent sequence of integrable simple functions  $s$ , we must show that the sequence  $\lambda n. \text{cond-exp } M F (s_n)$  is also convergent. Furthermore, we must show that this limit satisfies the properties of a conditional expectation. Unfortunately, we will only be able to show that this sequence converges in the L1-norm. Luckily, this is enough to show that the operator *cond-exp*  $M F$  preserves limits as a function from L1 to L1.

In anticipation of this result, we show that the conditional expectation operator is a contraction for simple functions. We first reformulate the lemma *real-cond-exp-abs*, which shows the statement for real-valued functions, using our definitions. Then we show the statement for simple functions via induction.

```
lemma cond-exp-contraction-real:
  fixes f :: 'a ⇒ real
  assumes integrable[measurable]: integrable M f
  shows AE x in M. norm (cond-exp M F f x) ≤ cond-exp M F (λx. norm (f x)) x
proof-
  have int: integrable M (λx. norm (f x)) using assms by blast
  have *: AE x in M. 0 ≤ cond-exp M F (λx. norm (f x)) x using cond-exp-real[THEN
AE-symmetric, OF integrable-norm[OF integrable]] real-cond-exp-ge-c[OF integrable-norm[OF
integrable], of 0] norm-ge-zero by fastforce
  have **: A ∈ sets F ⇒ (∫ x∈A. |f x| ∂M) = (∫ x∈A. real-cond-exp M F (λx.
norm (f x)) x ∂M) for A unfolding real-norm-def using assms integrable-abs
real-cond-exp-intA by blast
have norm-int: A ∈ sets F ⇒ (∫ x∈A. |f x| ∂M) = (∫ +x∈A. |f x| ∂M) for A
```

```

using assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce)
(meson subalg subalgebra-def subsetD)

have AE x in M. real-cond-exp M F (λx. norm (f x)) x ≥ 0 using int real-cond-exp-ge-c
by force
hence cond-exp-norm-int: A ∈ sets F ==> (ʃ x∈A. real-cond-exp M F (λx. norm
(f x)) x ∂M) = (ʃ+ x∈A. real-cond-exp M F (λx. norm (f x)) x ∂M) for A using
assms by (intro nn-set-integral-eq-set-integral[symmetric], blast, fastforce) (meson
subalg subalgebra-def subsetD)

have A ∈ sets F ==> (ʃ+ x∈A. |f x| ∂M) = (ʃ+ x∈A. real-cond-exp M F (λx.
norm (f x)) x ∂M) for A using ** norm-int cond-exp-norm-int by (auto simp
add: nn-integral-set-ennreal)
moreover have (λx. ennreal |f x|) ∈ borel-measurable M by measurable
moreover have (λx. ennreal (real-cond-exp M F (λx. norm (f x)) x)) ∈ borel-measurable
F by measurable
ultimately have AE x in M. nn-cond-exp M F (λx. ennreal |f x|) x = real-cond-exp
M F (λx. norm (f x)) x by (intro nn-cond-exp-charact[THEN AE-symmetric],
auto)
hence AE x in M. nn-cond-exp M F (λx. ennreal |f x|) x ≤ cond-exp M F (λx.
norm (f x)) x using cond-exp-real[OF int] by force
moreover have AE x in M. |real-cond-exp M F f x| = norm (cond-exp M F f x)
unfolding real-norm-def using cond-exp-real[OF assms] * by force
ultimately have AE x in M. ennreal (norm (cond-exp M F f x)) ≤ cond-exp M F
(λx. norm (f x)) x using real-cond-exp-abs[OF assms[THEN borel-measurable-integrable]]
by fastforce
hence AE x in M. enn2real (ennreal (norm (cond-exp M F f x))) ≤ enn2real
(cond-exp M F (λx. norm (f x)) x) using ennreal-le-iff2 by force
thus ?thesis using * by fastforce
qed

lemma cond-exp-contraction-simple:
fixes f :: 'a ⇒ 'b:: {second-countable-topology, banach}
assumes simple-function M f emeasure M {y ∈ space M. f y ≠ 0} ≠ ∞
shows AE x in M. norm (cond-exp M F f x) ≤ cond-exp M F (λx. norm (f x)) x
using assms
proof (induction rule: integrable-simple-function-induct)
case (cong f g)
hence ae: AE x in M. f x = g x by blast
hence AE x in M. cond-exp M F f x = cond-exp M F g x using cong has-cond-exp-simple
by (subst cond-exp-cong-AE) (auto intro!: has-cond-expD(2))
hence AE x in M. norm (cond-exp M F f x) = norm (cond-exp M F g x) by
force
moreover have AE x in M. cond-exp M F (λx. norm (f x)) x = cond-exp M F
(λx. norm (g x)) x using ae cong has-cond-exp-simple by (subst cond-exp-cong-AE)
(auto dest: has-cond-expD)
ultimately show ?case using cong(6) by fastforce
next
case (indicator A y)

```

**hence**  $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda a. \text{ indicator } A a *_R y) x = \text{cond-exp } M F (\text{indicator } A) x *_R y$  **by** blast  
**hence**  $\text{*}: \text{AE } x \text{ in } M. \text{ norm } (\text{cond-exp } M F (\lambda a. \text{ indicat-real } A a *_R y) x) \leq \text{norm } y$   
 $* \text{ cond-exp } M F (\lambda x. \text{ norm } (\text{indicat-real } A x)) x$  **using**  $\text{cond-exp-contraction-real}[\text{OF integrable-real-indicator, OF indicator}]$  **by** fastforce

**have**  $\text{AE } x \text{ in } M. \text{ norm } y * \text{ cond-exp } M F (\lambda x. \text{ norm } (\text{indicat-real } A x)) x = \text{norm } y * \text{ real-cond-exp } M F (\lambda x. \text{ norm } (\text{indicat-real } A x)) x$  **using**  $\text{cond-exp-real}[\text{OF integrable-real-indicator, OF indicator}]$  **by** fastforce

**moreover have**  $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. \text{ norm } y * \text{ norm } (\text{indicat-real } A x)) x = \text{real-cond-exp } M F (\lambda x. \text{ norm } y * \text{ norm } (\text{indicat-real } A x)) x$  **using**  $\text{indicator}$  **by** (intro cond-exp-real, auto)

**ultimately have**  $\text{AE } x \text{ in } M. \text{ norm } y * \text{ cond-exp } M F (\lambda x. \text{ norm } (\text{indicat-real } A x)) x = \text{cond-exp } M F (\lambda x. \text{ norm } y * \text{ norm } (\text{indicat-real } A x)) x$  **using**  $\text{real-cond-exp-cmult}[\text{of } \lambda x. \text{ norm } (\text{indicat-real } A x) \text{ norm } y]$  **indicator** **by** fastforce

**moreover have**  $(\lambda x. \text{ norm } y * \text{ norm } (\text{indicat-real } A x)) = (\lambda x. \text{ norm } (\text{indicat-real } A x *_R y))$  **by** force

**ultimately show** ?case **using** \* **by** force

**next**

**case** (add u v)

**have**  $\text{AE } x \text{ in } M. \text{ norm } (\text{cond-exp } M F (\lambda a. u a + v a) x) = \text{norm } (\text{cond-exp } M F u x + \text{cond-exp } M F v x)$  **using**  $\text{has-cond-exp-charact}(2)[\text{OF has-cond-exp-add, OF has-cond-exp-simple}(1,1), \text{ OF add}(1,2,3,4)]$  **by** fastforce

**moreover have**  $\text{AE } x \text{ in } M. \text{ norm } (\text{cond-exp } M F u x + \text{cond-exp } M F v x) \leq \text{norm } (\text{cond-exp } M F u x) + \text{norm } (\text{cond-exp } M F v x)$  **using**  $\text{norm-triangle-ineq}$  **by** blast

**moreover have**  $\text{AE } x \text{ in } M. \text{ norm } (\text{cond-exp } M F u x) + \text{norm } (\text{cond-exp } M F v x) \leq \text{cond-exp } M F (\lambda x. \text{ norm } (u x)) x + \text{cond-exp } M F (\lambda x. \text{ norm } (v x)) x$  **using** add(6,7) **by** fastforce

**moreover have**  $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. \text{ norm } (u x)) x + \text{cond-exp } M F (\lambda x. \text{ norm } (v x)) x = \text{cond-exp } M F (\lambda x. \text{ norm } (u x) + \text{norm } (v x)) x$  **using** integrable-simple-function[ $\text{OF add}(1,2)$ ] integrable-simple-function[ $\text{OF add}(3,4)$ ] **by** (intro has-cond-exp-charact(2)[ $\text{OF has-cond-exp-add}[\text{OF has-cond-exp-charact}(1,1)], \text{ THEN AE-symmetric}], auto intro: has-cond-exp-real)$

**moreover have**  $\text{AE } x \text{ in } M. \text{ cond-exp } M F (\lambda x. \text{ norm } (u x) + \text{norm } (v x)) x = \text{cond-exp } M F (\lambda x. \text{ norm } (u x + v x)) x$  **using** add(5) integrable-simple-function[ $\text{OF add}(1,2)$ ] integrable-simple-function[ $\text{OF add}(3,4)$ ] **by** (intro cond-exp-cong, auto)

**ultimately show** ?case **by** force

**qed**

**lemma** has-cond-exp-simple-lim:

**fixes**  $f :: 'a \Rightarrow 'b :: \{\text{second-countable-topology, banach}\}$

**assumes** integrable[measurable]: integrable  $M f$

**and**  $\bigwedge i. \text{ simple-function } M (s i)$

**and**  $\bigwedge i. \text{ emeasure } M \{y \in \text{space } M. s i y \neq 0\} \neq \infty$

**and**  $\bigwedge x. x \in \text{space } M \implies (\lambda i. s i x) \xrightarrow{} f x$

**and**  $\bigwedge x i. x \in \text{space } M \implies \text{norm } (s i x) \leq 2 * \text{norm } (f x)$

**obtains**  $r$

**where** strict-mono  $r$  has-cond-exp  $M F f (\lambda x. \text{ lim } (\lambda i. \text{ cond-exp } M F (s (r i)))$

$x))$   
 $\text{AE } x \text{ in } M. \text{ convergent } (\lambda i. \text{ cond-exp } M F (s (r i)) x)$   
**proof –**  
**have** [measurable]:  $(s i) \in \text{borel-measurable } M$  **for**  $i$  **using** assms(2) **by** (simp add: borel-measurable-simple-function)  
**have** integrable-s: integrable  $M (\lambda x. s i x)$  **for**  $i$  **using** assms integrable-simple-function by blast  
**have** integrable-4f: integrable  $M (\lambda x. 4 * \text{norm} (f x))$  **using** assms(1) **by** simp  
**have** integrable-2f: integrable  $M (\lambda x. 2 * \text{norm} (f x))$  **using** assms(1) **by** simp  
**have** integrable-2-cond-exp-norm-f: integrable  $M (\lambda x. 2 * \text{cond-exp } M F (\lambda x. \text{norm} (f x)) x)$  **by** fast  
  
**have** emeasure  $M \{y \in \text{space } M. s i y - s j y \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} + \text{emeasure } M \{y \in \text{space } M. s j y \neq 0\}$  **for**  $i j$  **using** simple-functionD(2)[OF assms(2)] **by** (intro order-trans[OF emeasure-mono emeasure-subadditive], auto)  
**hence** fin-sup: emeasure  $M \{y \in \text{space } M. s i y - s j y \neq 0\} \neq \infty$  **for**  $i j$  **using** assms(3) **by** (metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def)  
  
**have** emeasure  $M \{y \in \text{space } M. \text{norm} (s i y - s j y) \neq 0\} \leq \text{emeasure } M \{y \in \text{space } M. s i y \neq 0\} + \text{emeasure } M \{y \in \text{space } M. s j y \neq 0\}$  **for**  $i j$  **using** simple-functionD(2)[OF assms(2)] **by** (intro order-trans[OF emeasure-mono emeasure-subadditive], auto)  
**hence** fin-sup-norm: emeasure  $M \{y \in \text{space } M. \text{norm} (s i y - s j y) \neq 0\} \neq \infty$  **for**  $i j$  **using** assms(3) **by** (metis (mono-tags) ennreal-add-eq-top linorder-not-less top.not-eq-extremum infinity-ennreal-def)  
  
**have** Cauchy: Cauchy  $(\lambda n. s n x)$  **if**  $x \in \text{space } M$  **for**  $x$  **using** assms(4) LIMSEQ-imp-Cauchy that **by** blast  
**hence** bounded-range-s: bounded  $(\text{range} (\lambda n. s n x))$  **if**  $x \in \text{space } M$  **for**  $x$  **using** that cauchy-imp-bounded **by** fast

Since the sequence  $\lambda n. s n x$  is Cauchy for almost all  $x$ , we know that the diameter tends to zero almost everywhere.

Dominated convergence tells us that the integral of the diameter also converges to zero.

**have** AE  $x$  in  $M. (\lambda n. \text{diameter} \{s i x \mid i. n \leq i\}) \xrightarrow{} 0$  **using** Cauchy cauchy-iff-diameter-tends-to-zero-and-bounded **by** fast  
**moreover have**  $(\lambda x. \text{diameter} \{s i x \mid i. n \leq i\}) \in \text{borel-measurable } M$  **for**  $n$  **using** bounded-range-s borel-measurable-diameter **by** measurable  
**moreover have** AE  $x$  in  $M. \text{norm} (\text{diameter} \{s i x \mid i. n \leq i\}) \leq 4 * \text{norm} (f x)$  **for**  $n$   
**proof –**  
{  
fix  $x$  **assume**  $x: x \in \text{space } M$   
**have** diameter  $\{s i x \mid i. n \leq i\} \leq 2 * \text{norm} (f x) + 2 * \text{norm} (f x)$  **by** (intro diameter-le, blast, subst dist-norm[symmetric], intro dist-triangle3[THEN

```

order-trans, of 0], intro add-mono) (auto intro: assms(5)[OF x])
  hence norm (diameter {s i x | i. n ≤ i}) ≤ 4 * norm (f x) using diameter-ge-0[OF bounded-subset[OF bounded-range-s], OF x, of {s i x | i. n ≤ i}] by force
  }
  thus ?thesis by fast
qed
ultimately have diameter-tends-to-zero: (λn. LINT x|M. diameter {s i x | i. n ≤ i}) —→ 0 by (intro integral-dominated-convergence[OF borel-measurable-const[of 0] - integrable-4f, simplified]) (fast+)

have diameter-integrable: integrable M (λx. diameter {s i x | i. n ≤ i}) for n using assms(1,5)
  by (intro integrable-bound-diameter[OF bounded-range-s integrable-2f], auto)

have dist-integrable: integrable M (λx. dist (s i x) (s j x)) for i j using assms(5)
dist-triangle3[of s i - - 0, THEN order-trans, OF add-mono, of - 2 * norm (f -)]
  by (intro Bochner-Integration.integrable-bound[OF integrable-4f]) fastforce+

```

Since  $\text{cond-exp } M F$  is a contraction for simple functions, the following sequence of integral values is also Cauchy.

This follows, since the distance between the terms of this sequence are always less than or equal to the diameter, which itself converges to zero.

Hence, we obtain a subsequence which is Cauchy almost everywhere.

```

have ∃ N. ∀ i≥N. ∀ j≥N. LINT x|M. norm (cond-exp M F (s i) x - cond-exp M F (s j) x) < e if e-pos: e > 0 for e
proof -
  obtain N where *: LINT x|M. diameter {s i x | i. n ≤ i} < e if n ≥ N for n using that order-tends-to-zero[THEN iffD1, OF diameter-tends-to-zero, unfolded eventually-sequentially] e-pos by presburger
  {
    fix i j x assume asm: i ≥ N j ≥ N x ∈ space M
    have case-prod dist ‘({s i x | i. N ≤ i} × {s i x | i. N ≤ i}) = case-prod (λi j. dist (s i x) (s j x)) ‘({N..} × {N..}) by fast
    hence diameter {s i x | i. N ≤ i} = (SUP (i, j) ∈ {N..} × {N..}. dist (s i x) (s j x)) unfolding diameter-def by auto
    moreover have (SUP (i, j) ∈ {N..} × {N..}. dist (s i x) (s j x)) ≥ dist (s i x) (s j x) using asm bounded-imp-bdd-above[OF bounded-imp-dist-bounded, OF bounded-range-s] by (intro cSup-upper, auto)
    ultimately have diameter {s i x | i. N ≤ i} ≥ dist (s i x) (s j x) by presburger
  }
  hence LINT x|M. dist (s i x) (s j x) < e if i ≥ N j ≥ N for i j using that * by (intro integral-mono[OF dist-integrable diameter-integrable, THEN order.strict-trans1], blast+)
  moreover have LINT x|M. norm (cond-exp M F (s i) x - cond-exp M F (s j) x) ≤ LINT x|M. dist (s i x) (s j x) for i j

```

**proof** –

**have**  $LINT x|M. \text{norm}(\text{cond-exp } M F (s i) x - \text{cond-exp } M F (s j) x) = LINT x|M. \text{norm}(\text{cond-exp } M F (s i) x + - 1 *_R \text{cond-exp } M F (s j) x)$  **unfolding dist-norm by simp**

**also have** ... =  $LINT x|M. \text{norm}(\text{cond-exp } M F (\lambda x. s i x - s j x) x)$  **using has-cond-exp-charact(2)[OF has-cond-exp-add[OF - has-cond-exp-scaleR-right, OF has-cond-exp-charact(1,1), OF has-cond-exp-simple(1,1)[OF assms(2,3)]], THEN AE-symmetric, of  $i - 1 j$ ] by (intro integral-cong-AE) force+**

**also have** ...  $\leq LINT x|M. \text{cond-exp } M F (\lambda x. \text{norm}(s i x - s j x)) x$  **using cond-exp-contraction-simple[OF - fin-sup, of  $i j$ ] integrable-cond-exp assms(2) by (intro integral-mono-AE, fast+)**

**also have** ... =  $LINT x|M. \text{norm}(s i x - s j x)$  **unfolding set-integral-space(1)[OF integrable-cond-exp, symmetric] set-integral-space[OF dist-integrable[unfolded dist-norm], symmetric] by (intro has-cond-expD(1)[OF has-cond-exp-simple[OF - fin-sup-norm], symmetric]) (metis assms(2) simple-function-compose1 simple-function-diff, metis sets.top subalg subalgebra-def)**

**finally show** ?thesis **unfolding dist-norm .**

**qed**

**ultimately show** ?thesis **using order.strict-trans1 by meson**

**qed**

**then obtain**  $r$  **where strict-mono-r: strict-mono  $r$  and AE-Cauchy: AE  $x$  in  $M$ . Cauchy ( $\lambda i. \text{cond-exp } M F (s (r i)) x$ )**

**by** (rule cauchy-L1-AE-cauchy-subseq[OF integrable-cond-exp], auto)

**hence** ae-lim-cond-exp: AE  $x$  in  $M$ . ( $\lambda n. \text{cond-exp } M F (s (r n)) x$ )  $\longrightarrow$  lim ( $\lambda n. \text{cond-exp } M F (s (r n)) x$ ) **using Cauchy-convergent-iff convergent-LIMSEQ-iff by fastforce**

Now that we have a candidate for the conditional expectation, we must show that it actually has the required properties.

Dominated convergence shows that this limit is indeed integrable.

Here, we again use the fact that conditional expectation is a contraction on simple functions.

**have** cond-exp-bounded: AE  $x$  in  $M$ .  $\text{norm}(\text{cond-exp } M F (s (r n)) x) \leq \text{cond-exp } M F (\lambda x. 2 * \text{norm}(f x)) x$  **for n**

**proof** –

**have** AE  $x$  in  $M$ .  $\text{norm}(\text{cond-exp } M F (s (r n)) x) \leq \text{cond-exp } M F (\lambda x. \text{norm}(s (r n) x)) x$  **by** (rule cond-exp-contraction-simple[OF assms(2,3)])

**moreover have** AE  $x$  in  $M$ . real-cond-exp  $M F (\lambda x. \text{norm}(s (r n) x)) x \leq \text{real-cond-exp } M F (\lambda x. 2 * \text{norm}(f x)) x$  **using integrable-s integrable-2f assms(5) by (intro real-cond-exp-mono, auto)**

**ultimately show** ?thesis **using cond-exp-real[OF integrable-norm, OF integrable-s, of  $r n$ ] cond-exp-real[OF integrable-2f] by force**

**qed**

**have** lim-integrable: integrable  $M$  ( $\lambda x. \lim(\lambda i. \text{cond-exp } M F (s (r i)) x)$ ) **by** (intro integrable-dominated-convergence[OF - borel-measurable-cond-exp' integrable-cond-exp ae-lim-cond-exp cond-exp-bounded], simp)

Moreover, we can use the DCT twice to show that the conditional expecta-

tion property holds, i.e. the value of the integral of the candidate, agrees with  $f$  on sets  $A \in \text{sets } F$ .

```
{
  fix A assume A-in-sets-F: A ∈ sets F
  have AE x in M. norm (indicator A x *R cond-exp M F (s (r n)) x) ≤ cond-exp
    M F (λx. 2 * norm (f x)) x for n
  proof -
    have AE x in M. norm (indicator A x *R cond-exp M F (s (r n)) x) ≤ norm
      (cond-exp M F (s (r n)) x) unfolding indicator-def by simp
      thus ?thesis using cond-exp-bounded[of n] by force
    qed
    hence lim-cond-exp-int: (λn. LINT x:A|M. cond-exp M F (s (r n)) x) —→
      (LINT x:A|M. lim (λn. cond-exp M F (s (r n)) x))
    using ae-lim-cond-exp measurable-from-subalg[OF subalg borel-measurable-indicator,
      OF A-in-sets-F] cond-exp-bounded
    unfolding set-lebesgue-integral-def
    by (intro integral-dominated-convergence[OF borel-measurable-scaleR borel-measurable-scaleR
      integrable-cond-exp]) (fastforce simp add: tendsto-scaleR)+

    have AE x in M. norm (indicator A x *R s (r n) x) ≤ 2 * norm (f x) for n
    proof -
      have AE x in M. norm (indicator A x *R s (r n) x) ≤ norm (s (r n) x)
      unfolding indicator-def by simp
      thus ?thesis using assms(5)[of - r n] by fastforce
    qed
    hence lim-s-int: (λn. LINT x:A|M. s (r n) x) —→ (LINT x:A|M. f x)
    using measurable-from-subalg[OF subalg borel-measurable-indicator, OF A-in-sets-F]
    LIMSEQ-subseq-LIMSEQ[OF assms(4) strict-mono-r] assms(5)
    unfolding set-lebesgue-integral-def comp-def
    by (intro integral-dominated-convergence[OF borel-measurable-scaleR borel-measurable-scaleR
      integrable-2f]) (fastforce simp add: tendsto-scaleR)+

    have (LINT x:A|M. lim (λn. cond-exp M F (s (r n)) x)) = lim (λn. LINT
      x:A|M. cond-exp M F (s (r n)) x) using limI[OF lim-cond-exp-int] by argo
    also have ... = lim (λn. LINT x:A|M. s (r n) x) using has-cond-expD(1)[OF
      has-cond-exp-simple[OF assms(2,3)] A-in-sets-F, symmetric] by presburger
    also have ... = (LINT x:A|M. f x) using limI[OF lim-s-int] by argo
    finally have (LINT x:A|M. lim (λn. cond-exp M F (s (r n)) x)) = (LINT
      x:A|M. f x) .
  }
}
```

Putting it all together, we have the statement we are looking for.

```

  hence has-cond-exp M F f (λx. lim (λi. cond-exp M F (s (r i)) x)) using
    assms(1) lim-integrable by (intro has-cond-expI', auto)
  thus thesis using AE-Cauchy Cauchy-convergent strict-mono-r by (auto intro!
    that)
  qed
}
```

Now, we can show that the conditional expectation is well-defined for all

integrable functions.

```

corollary has-cond-expI:
  fixes f :: 'a ⇒ 'b:{second-countable-topology,banach}
  assumes integrable M f
  shows has-cond-exp M F f (cond-exp M F f)
proof –
  obtain s where s-is: ⋀ i. simple-function M (s i) ⋀ i. emeasure M {y ∈ space M.
  s i y ≠ 0} ≠ ∞ ⋀ x. x ∈ space M ⇒ (λ i. s i x) —→ fx ⋀ x. x ∈ space M ⇒
  norm (s i x) ≤ 2 * norm (fx) using integrable-implies-simple-function-sequence[OF
  assms] by blast
  show ?thesis using has-cond-exp-simple-lim[OF assms s-is] has-cond-exp-charact(1)
  by metis
qed
```

## 2.2 Properties

The defining property of the conditional expectation now always holds, given that the function  $f$  is integrable.

```

lemma cond-exp-set-integral:
  fixes f :: 'a ⇒ 'b:{second-countable-topology,banach}
  assumes integrable M f A ∈ sets F
  shows (∫ x ∈ A. f x ∂M) = (∫ x ∈ A. cond-exp M F f x ∂M)
  using has-cond-expD(1)[OF has-cond-expI, OF assms] by argo
```

The following property of the conditional expectation is called the "Tower Property".

```

lemma cond-exp-nested-subalg:
  fixes f :: 'a ⇒ 'b:{second-countable-topology,banach}
  assumes integrable M f subalgebra M G subalgebra G F
  shows AE ξ in M. cond-exp M F f ξ = cond-exp M F (cond-exp M G f) ξ
  using has-cond-expI assms sigma-finite-subalgebra-def by (auto intro!: has-cond-exp-nested-subalg[THEN
  has-cond-exp-charact(2), THEN AE-symmetric] sigma-finite-subalgebra.has-cond-expI[OF
  sigma-finite-subalgebra.intro[OF assms(2)]] nested-subalg-is-sigma-finite)
```

The conditional expectation is linear.

```

lemma cond-exp-add:
  fixes f :: 'a ⇒ 'b:{second-countable-topology,banach}
  assumes integrable M f integrable M g
  shows AE x in M. cond-exp M F (λx. f x + g x) x = cond-exp M F f x +
  cond-exp M F g x
  using has-cond-exp-add[OF has-cond-expI(1,1), OF assms, THEN has-cond-exp-charact(2)]
  .
```

```

lemma cond-exp-diff:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach}
  assumes integrable M f integrable M g
  shows AE x in M. cond-exp M F (λx. f x - g x) x = cond-exp M F f x -
  cond-exp M F g x
```

**using** *has-cond-exp-add*[*OF - has-cond-exp-scaleR-right, OF has-cond-expI(1,1), OF assms, THEN has-cond-exp-charact(2), of -1*] **by** *simp*

```

lemma cond-exp-diff':
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach}
  assumes integrable M f integrable M g
  shows AE x in M. cond-exp M F (f - g) x = cond-exp M F f x - cond-exp M F g x
  unfolding fun-diff-def using assms by (rule cond-exp-diff)
lemma cond-exp-scaleR-left:
  fixes f :: 'a  $\Rightarrow$  real
  assumes integrable M f
  shows AE x in M. cond-exp M F ( $\lambda x. f x *_R c$ ) x = cond-exp M F f x *_R c
  using cond-exp-set-integral[OF assms] subalg assms unfolding subalgebra-def
  by (intro cond-exp-charact,
    subst set-integral-scaleR-left, blast, intro assms,
    subst set-integral-scaleR-left, blast, intro integrable-cond-exp)
  auto

```

The conditional expectation operator is a contraction, i.e. a bounded linear operator with operator norm less than or equal to 1.

To show this we first obtain a subsequence  $\lambda x. i. s(r i) x$ , such that  $\lambda i. cond-exp M F (s(r i)) x$  converges to  $cond-exp M F f x$  a.e. Afterwards, we obtain a sub-subsequence  $\lambda x. i. s(r(r' i)) x$ , such that  $\lambda i. cond-exp M F (\lambda x. norm(s(r i))) x$  converges to  $cond-exp M F (\lambda x. norm(f x)) x$  a.e. Finally, we show that the inequality holds by showing that the terms of the subsequences obey the inequality and the fact that a subsequence of a convergent sequence converges to the same limit.

```

lemma cond-exp-contraction:
  fixes f :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach}
  assumes integrable M f
  shows AE x in M. norm (cond-exp M F f x)  $\leq$  cond-exp M F ( $\lambda x. norm(f x)$ )
  x
proof -
  obtain s where s:  $\bigwedge i. simple-function M (s i) \bigwedge i. emeasure M \{y \in space M. s i y \neq 0\} \neq \infty \bigwedge x. x \in space M \implies (\lambda i. s i x) \longrightarrow f x \bigwedge i. x \in space M \implies norm(s i x) \leq 2 * norm(f x)$ 
  by (blast intro: integrable-implies-simple-function-sequence[OF assms])
  obtain r where r: strict-mono r and has-cond-exp M F f ( $\lambda x. lim(\lambda i. cond-exp M F (s(r i)) x)$ ) AE x in M. ( $\lambda i. cond-exp M F (s(r i)) x$ )  $\longrightarrow$  lim ( $\lambda i. cond-exp M F (s(r i)) x$ )
  using has-cond-exp-simple-lim[OF assms s] unfolding convergent-LIMSEQ-iff
  by blast
  hence r-tendsto: AE x in M. ( $\lambda i. cond-exp M F (s(r i)) x$ )  $\longrightarrow$  cond-exp M F f x using has-cond-exp-charact(2) by force

```

```

have norm-s-r:  $\bigwedge i. \text{simple-function } M (\lambda x. \text{norm} (s (r i) x)) \bigwedge i. \text{emeasure } M$   

 $\{y \in \text{space } M. \text{norm} (s (r i) y) \neq 0\} \neq \emptyset \bigwedge x. x \in \text{space } M \implies (\lambda i. \text{norm} (s (r i) x)) \longrightarrow \text{norm} (f x) \bigwedge i. x. x \in \text{space } M \implies \text{norm} (\text{norm} (s (r i) x)) \leq 2 * \text{norm} (\text{norm} (f x))$   

using s by (auto intro: LIMSEQ-subseq-LIMSEQ[OF tendsto-norm r, unfolded comp-def] simple-function-compose1)  

  

obtain r' where r': strict-mono r' and has-cond-exp M F ( $\lambda x. \text{norm} (f x)$ ) ( $\lambda x.$   

 $\text{lim} (\lambda i. \text{cond-exp } M F (\lambda x. \text{norm} (s (r (r' i)) x))) x$ ) AE x in M. ( $\lambda i. \text{cond-exp } M F (\lambda x. \text{norm} (s (r (r' i)) x)) x$ )  $\longrightarrow \text{lim} (\lambda i. \text{cond-exp } M F (\lambda x. \text{norm} (s (r (r' i)) x)) x)$  using has-cond-exp-simple-lim[OF integrable-norm norm-s-r, OF assms] unfolding convergent-LIMSEQ-iff by blast  

hence r'-tendsto: AE x in M. ( $\lambda i. \text{cond-exp } M F (\lambda x. \text{norm} (s (r (r' i)) x)) x$ )  $\longrightarrow \text{cond-exp } M F (\lambda x. \text{norm} (f x)) x$  using has-cond-exp-charact(2) by force  

  

have AE x in M.  $\forall i. \text{norm} (\text{cond-exp } M F (s (r (r' i)))) x \leq \text{cond-exp } M F (\lambda x.$   

 $\text{norm} (s (r (r' i)) x)) x$  using s by (auto intro: cond-exp-contraction-simple simp add: AE-all-countable)  

moreover have AE x in M. ( $\lambda i. \text{norm} (\text{cond-exp } M F (s (r (r' i)))) x$ )  $\longrightarrow$   

 $\text{norm} (\text{cond-exp } M F f x)$  using r-tendsto LIMSEQ-subseq-LIMSEQ[OF tendsto-norm r', unfolded comp-def] by fast  

ultimately show ?thesis using LIMSEQ-le r'-tendsto by fast  

qed

```

The following lemmas are called "pulling out whats known". We first show the statement for real-valued functions using the lemma *real-cond-exp-intg*, which is already present. We then show it for arbitrary *g* using the lecture notes of Gordan Zitkovic for the course "Theory of Probability I" [4].

```

lemma cond-exp-measurable-mult:  

fixes f g :: 'a  $\Rightarrow$  real  

assumes [measurable]: integrable M ( $\lambda x. f x * g x$ ) integrable M g f  $\in$  borel-measurable F  

shows integrable M ( $\lambda x. f x * \text{cond-exp } M F g x$ )  

AE x in M. cond-exp M F ( $\lambda x. f x * g x$ ) x = f x * cond-exp M F g x  

proof –  

show integrable: integrable M ( $\lambda x. f x * \text{cond-exp } M F g x$ ) using cond-exp-real[OF assms(2)] by (intro integrable-cong-AE-imp[OF real-cond-exp-intg(1), OF assms(1,3)] assms(2)[THEN borel-measurable-integrable]) measurable-from-subalg[OF subalg]  

auto  

interpret sigma-finite-measure restr-to-subalg M F by (rule sigma-fin-subalg)  

{  

fix A assume asm: A  $\in$  sets F  

hence asm': A  $\in$  sets M using subalg by (fastforce simp add: subalgebra-def)  

have set-lebesgue-integral M A ( $\text{cond-exp } M F (\lambda x. f x * g x)$ ) = set-lebesgue-integral M A ( $\lambda x. f x * g x$ ) by (simp add: cond-exp-set-integral[OF assms(1) asm])  

also have ... = set-lebesgue-integral M A ( $\lambda x. f x * \text{real-cond-exp } M F g$   

x) using borel-measurable-times[OF borel-measurable-indicator[OF asm] assms(3)]  

borel-measurable-integrable[OF assms(2)] integrable-mult-indicator[OF asm' assms(1)]  

by (fastforce simp add: set-lebesgue-integral-def mult.assoc[symmetric] intro: real-cond-exp-intg(2)[symmetric])

```

**also have** ... = set-lebesgue-integral M A ( $\lambda x. f x * \text{cond-exp } M F g x$ ) **using**  
 cond-exp-real[*OF assms(2)*] *asm'* borel-measurable-cond-exp' borel-measurable-cond-exp2  
 measurable-from-subalg[*OF subalg assms(3)*] **by** (auto simp add: set-lebesgue-integral-def  
 intro: integral-cong-AE)  
**finally have** set-lebesgue-integral M A ( $\text{cond-exp } M F (\lambda x. f x * g x) = (\int x \in A.$   
 $f x * \text{cond-exp } M F g x) \partial M$ ) .  
}

**hence** AE x in restr-to-subalg M F. cond-exp M F ( $\lambda x. f x * g x$ ) x = f  
 $x * \text{cond-exp } M F g x$  **by** (intro density-unique-banach integrable-cond-exp integrable  
 integrable-in-subalg subalg, measurable, simp add: set-lebesgue-integral-def  
 integral-subalgebra2[*OF subalg*] sets-restr-to-subalg[*OF subalg*])

**thus** AE x in M. cond-exp M F ( $\lambda x. f x * g x$ ) x = f x \* cond-exp M F g x **by**  
 (rule AE-restr-to-subalg[*OF subalg*])

**qed**

**lemma** cond-exp-measurable-scaleR:  
**fixes** f :: 'a  $\Rightarrow$  real **and** g :: 'a  $\Rightarrow$  'b :: {second-countable-topology, banach}  
**assumes** [measurable]: integrable M ( $\lambda x. f x *_R g x$ ) integrable M g f  $\in$  borel-measurable  
 F  
**shows** integrable M ( $\lambda x. f x *_R \text{cond-exp } M F g x$ )  
 $\text{AE } x \text{ in } M. \text{cond-exp } M F (\lambda x. f x *_R g x) x = f x *_R \text{cond-exp } M F g x$   
**proof** –  
 let ?F = restr-to-subalg M F  
**have** subalg': subalgebra M (restr-to-subalg M F) **by** (metis sets-eq-imp-space-eq  
 sets-restr-to-subalg subalg subalgebra-def)  
{
 fix z **assume** asm[measurable]: integrable M ( $\lambda x. z x *_R g x$ ) z  $\in$  borel-measurable  
 ?F  
**hence** asm'[measurable]: z  $\in$  borel-measurable F **using** measurable-in-subalg'  
 subalg **by** blast  
**have** integrable M ( $\lambda x. z x *_R \text{cond-exp } M F g x$ ) LINT x|M. z x \*\_R g x =  
 LINT x|M. z x \*\_R cond-exp M F g x  
**proof** –  
 obtain s where s-is:  $\bigwedge i. \text{simple-function } ?F (s i) \wedge x \in \text{space } ?F \implies (\lambda i. s i x) \longrightarrow z x \bigwedge i. x \in \text{space } ?F \implies \text{norm } (s i x) \leq 2 * \text{norm } (z x)$  **using**  
 borel-measurable-implies-sequence-metric[*OF asm(2), of 0*] **by** force

We need to apply the dominated convergence theorem twice, therefore we need to show the following prerequisites.

**have** s-scaleR-g-tendsto: AE x in M. ( $\lambda i. s i x *_R g x \longrightarrow z x *_R g x$ )  $\longrightarrow$  z x \*\_R g x  
**using** s-is(2) **by** (simp add: space-restr-to-subalg tendsto-scaleR)  
**have** s-scaleR-cond-exp-g-tendsto: AE x in ?F. ( $\lambda i. s i x *_R \text{cond-exp } M F g x \longrightarrow z x *_R \text{cond-exp } M F g x$ )  $\longrightarrow$  z x \*\_R cond-exp M F g x **using** s-is(2) **by** (simp add: tendsto-scaleR)

**have** s-scaleR-g-meas: ( $\lambda x. s i x *_R g x$ )  $\in$  borel-measurable M **for** i **using**  
 s-is(1)[THEN borel-measurable-simple-function, THEN subalg'[THEN measurable-from-subalg]] **by** simp  
**have** s-scaleR-cond-exp-g-meas: ( $\lambda x. s i x *_R \text{cond-exp } M F g x$ )  $\in$  borel-measurable  
 ?F **for** i **using** s-is(1)[THEN borel-measurable-simple-function] measurable-in-subalg[*OF*

*subalg borel-measurable-cond-exp] by (fastforce intro: borel-measurable-scaleR)*

```

have s-scaleR-g-AE-bdd: AE x in M. norm (s i x *R g x) ≤ 2 * norm
(z x *R g x) for i using s-is(3) by (fastforce simp add: space-restr-to-subalg
mult.assoc[symmetric] mult-right-mono)
{
  fix i
  have asm: integrable M (λx. norm (z x) * norm (g x)) using asm(1)[THEN
integrable-norm] by simp
  have AE x in ?F. norm (s i x *R cond-exp M F g x) ≤ 2 * norm (z x) *
norm (cond-exp M F g x) using s-is(3) by (fastforce simp add: mult-mono)
  moreover have AE x in ?F. norm (z x) * cond-exp M F (λx. norm (g x)) x =
cond-exp M F (λx. norm (z x) * norm (g x)) x by (rule cond-exp-measurable-mult(2)[THEN
AE-symmetric, OF asm integrable-norm, OF assms(2), THEN AE-restr-to-subalg2[OF
subalg]], auto)
  ultimately have AE x in ?F. norm (s i x *R cond-exp M F g x) ≤ 2 *
cond-exp M F (λx. norm (z x *R g x)) x using cond-exp-contraction[OF assms(2),
THEN AE-restr-to-subalg2[OF subalg]] order-trans[OF - mult-mono] by fastforce
}
note s-scaleR-cond-exp-g-AE-bdd = this

```

In the following section we need to pay attention to which measures we are using for integration. The rhs is F-measurable while the lhs is only M-measurable.

```

{
  fix i
  have s-meas-M[measurable]: s i ∈ borel-measurable M by (meson borel-measurable-simple-function
measurable-from-subalg s-is(1) subalg')
  have s-meas-F[measurable]: s i ∈ borel-measurable F by (meson borel-measurable-simple-function
measurable-in-subalg' s-is(1) subalg')

  have s-scaleR-eq: s i x *R h x = (∑ y ∈ s i ` space M. (indicator (s i - ` {y} *
  ∩ space M) x *R y) *R h x) if x ∈ space M for x and h :: 'a ⇒ 'b
    using simple-function-indicator-representation-banach[OF s-is(1), of x
i] that unfolding space-restr-to-subalg scaleR-left.sum[of - - h x, symmetric] by
presburger

  have LINT x|M. s i x *R g x = LINT x|M. (∑ y ∈ s i ` space M. indicator (s i - ` {y} *
  ∩ space M) x *R y *R g x) using s-scaleR-eq by (intro
Bochner-Integration.integral-cong) auto
  also have ... = (∑ y ∈ s i ` space M. LINT x|M. indicator (s i - ` {y} *
  ∩ space M) x *R y *R g x) by (intro Bochner-Integration.integral-sum
integrable-mult-indicator[OF - integrable-scaleR-right] assms(2)) simp
  also have ... = (∑ y ∈ s i ` space M. y *R set-lebesgue-integral M (s i - ` {y} *
  ∩ space M) g) by (simp only: set-lebesgue-integral-def[symmetric]) simp
  also have ... = (∑ y ∈ s i ` space M. y *R set-lebesgue-integral M (s i - ` {y} *
  ∩ space M) (cond-exp M F g)) using assms(2) subalg borel-measurable-vimage[OF
s-meas-F] by (subst cond-exp-set-integral, auto simp add: subalgebra-def)
  also have ... = (∑ y ∈ s i ` space M. LINT x|M. indicator (s i - ` {y} ∩ space

```

```

 $M) x *_R y *_R \text{cond-exp } M F g x$  by (simp only: set-lebesgue-integral-def[symmetric])
simp
also have ... =  $\text{LINT } x|M. (\sum_{y \in s} i \cdot \text{space } M. \text{indicator } (s i - \{y\}) \cap \text{space } M) x *_R y *_R \text{cond-exp } M F g x$  by (intro Bochner-Integration.integral-sum[symmetric]
integrable-mult-indicator[OF - integrable-scaleR-right]) auto
also have ... =  $\text{LINT } x|M. s i x *_R \text{cond-exp } M F g x$  using s-scaleR-eq
by (intro Bochner-Integration.integral-cong) auto
finally have  $\text{LINT } x|M. s i x *_R g x = \text{LINT } x|?F. s i x *_R \text{cond-exp } M F$ 
 $g x$  by (simp add: integral-subalgebra2[OF subalg])
}
note integral-s-eq = this

```

Now we just plug in the results we obtained into DCT, and use the fact that limits are unique.

```

show integrable  $M (\lambda x. z x *_R \text{cond-exp } M F g x)$  using s-scaleR-cond-exp-g-meas
asm(2) borel-measurable-cond-exp' by (intro integrable-from-subalg[OF subalg] inte-
grable-cond-exp integrable-dominated-convergence[OF --- s-scaleR-cond-exp-g-tendsto
s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] inte-
grable-in-subalg measurable-in-subalg subalg)
}

have  $(\lambda i. \text{LINT } x|M. s i x *_R g x) \longrightarrow \text{LINT } x|M. z x *_R g x$  using
s-scaleR-g-meas asm(1)[THEN integrable-norm] asm' borel-measurable-cond-exp'
by (intro integral-dominated-convergence[OF --- s-scaleR-g-tendsto s-scaleR-g-AE-bdd])
(auto intro: measurable-from-subalg[OF subalg])
moreover have  $(\lambda i. \text{LINT } x|?F. s i x *_R \text{cond-exp } M F g x) \longrightarrow$ 
 $\text{LINT } x|?F. z x *_R \text{cond-exp } M F g x$  using s-scaleR-cond-exp-g-meas asm(2)
borel-measurable-cond-exp' by (intro integral-dominated-convergence[OF --- s-scaleR-cond-exp-g-tendsto
s-scaleR-cond-exp-g-AE-bdd]) (auto intro: measurable-from-subalg[OF subalg] inte-
grable-in-subalg measurable-in-subalg subalg)
ultimately show  $\text{LINT } x|M. z x *_R g x = \text{LINT } x|M. z x *_R \text{cond-exp }$ 
 $M F g x$  using integral-s-eq using subalg by (simp add: LIMSEQ-unique inte-
gral-subalgebra2)
qed
}
note * = this

```

The main statement now follows with  $z = (\lambda x. \text{indicat-real } A x * f x)$ .

```

show integrable  $M (\lambda x. f x *_R \text{cond-exp } M F g x)$  using * assms measurable-in-subalg[OF subalg] by blast

{
  fix  $A$  assume asm: A ∈ F
  hence integrable  $M (\lambda x. \text{indicat-real } A x *_R f x *_R g x)$  using subalg by
  (fastforce simp add: subalgebra-def intro!: integrable-mult-indicator assms(1))
  hence set-lebesgue-integral  $M A (\lambda x. f x *_R g x) = \text{set-lebesgue-integral } M A$ 
   $(\lambda x. f x *_R \text{cond-exp } M F g x)$  unfolding set-lebesgue-integral-def using asm by
  (auto intro!: * measurable-in-subalg[OF subalg])
}
thus AE x in M. cond-exp M F (λx. f x *_R g x) x = f x *_R cond-exp M F g x

```

```

using borel-measurable-cond-exp by (intro cond-exp-charact, auto intro!: * assms
measurable-in-subalg[OF subalg])
qed

lemma cond-exp-sum [intro, simp]:
  fixes f :: 't ⇒ 'a ⇒ 'b :: {second-countable-topology, banach}
  assumes [measurable]: ∀ i. integrable M (f i)
  shows AE x in M. cond-exp M F (λx. ∑ i∈I. f i x) x = (∑ i∈I. cond-exp M F
(f i) x)
proof (rule has-cond-exp-charact, intro has-cond-expI')
  fix A assume [measurable]: A ∈ sets F
  then have A-meas [measurable]: A ∈ sets M by (meson subsetD subalg subalgebra-def)

  have (∫ x∈A. (∑ i∈I. f i x)∂M) = (∫ x. (∑ i∈I. indicator A x ∗_R f i x)∂M)
  unfolding set-lebesgue-integral-def by (simp add: scaleR-sum-right)
  also have ... = (∑ i∈I. (∫ x. indicator A x ∗_R f i x ∂M)) using assms by (auto
intro!: Bochner-Integration.integral-sum integrable-mult-indicator)
  also have ... = (∑ i∈I. (∫ x. indicator A x ∗_R cond-exp M F (f i) x ∂M)) using
cond-exp-set-integral[OF assms] by (simp add: set-lebesgue-integral-def)
  also have ... = (∫ x. (∑ i∈I. indicator A x ∗_R cond-exp M F (f i) x)∂M)
  using assms by (auto intro!: Bochner-Integration.integral-sum[symmetric] inte-
grable-mult-indicator)
  also have ... = (∫ x∈A. (∑ i∈I. cond-exp M F (f i) x)∂M) unfolding set-lebesgue-integral-def
by (simp add: scaleR-sum-right)
  finally show (∫ x∈A. (∑ i∈I. f i x)∂M) = (∫ x∈A. (∑ i∈I. cond-exp M F (f i)
x)∂M) by auto
qed (auto simp add: assms integrable-cond-exp)

```

## 2.3 Linearly Ordered Banach Spaces

In this subsection we show monotonicity results concerning the conditional expectation operator.

```

lemma cond-exp-gr-c:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
  assumes integrable M f AE x in M. f x > c
  shows AE x in M. cond-exp M F f x > c
proof -
  define X where X = {x ∈ space M. cond-exp M F f x ≤ c}
  have [measurable]: X ∈ sets F unfolding X-def by measurable (metis sets.top
subalg subalgebra-def)
  hence X-in-M: X ∈ sets M using sets-restr-to-subalg subalg subalgebra-def by
blast
  have emeasure M X = 0
  proof (rule ccontr)
    assume emeasure M X ≠ 0
    have emeasure (restr-to-subalg M F) X = emeasure M X by (simp add: eme-
sure-restr-to-subalg subalg)
  
```

```

hence emeasure (restr-to-subalg M F) X > 0 using ⊢(emeasure M X) = 0
gr-zeroI by auto
then obtain A where A: A ∈ sets (restr-to-subalg M F) A ⊆ X emeasure
(restr-to-subalg M F) A > 0 emeasure (restr-to-subalg M F) A < ∞
using sigma-fin-subalg by (metis emeasure-notin-sets ennreal-0 infinity-ennreal-def
le-less-linear neq-top-trans not-gr-zero order-refl sigma-finite-measure.approx-PInf-emeasure-with-finite)
hence [simp]: A ∈ sets F using subalg sets-restr-to-subalg by blast
hence A-in-sets-M[simp]: A ∈ sets M using sets-restr-to-subalg subalg subal-
gebra-def by blast
have [simp]: set-integrable M A (λx. c) using A subalg by (auto simp add:
set-integrable-def emeasure-restr-to-subalg)
have [simp]: set-integrable M A f unfolding set-integrable-def by (rule inte-
grable-mult-indicator, auto simp add: assms(1))
have AE x in M. indicator A x *R c = indicator A x *R f x
proof (rule integral-eq-mono-AE-eq-AE)
have (ʃ x∈A. c ∂M) ≤ (ʃ x∈A. f x ∂M) using assms(2) by (intro set-integral-mono-AE-banach)
auto
moreover
{
  have (ʃ x∈A. f x ∂M) = (ʃ x∈A. cond-exp M F f x ∂M) by (rule
cond-exp-set-integral, auto simp add: assms)
also have ... ≤ (ʃ x∈A. c ∂M) using A by (auto intro!: set-integral-mono-banach
simp add: X-def)
finally have (ʃ x∈A. f x ∂M) ≤ (ʃ x∈A. c ∂M) by simp
}
ultimately show LINT x|M. indicator A x *R c = LINT x|M. indicator A
x *R f x unfolding set-lebesgue-integral-def by simp
show AE x in M. indicator A x *R c ≤ indicator A x *R f x using assms by
(auto simp add: X-def indicator-def)
qed (auto simp add: set-integrable-def[symmetric])
hence AE x∈A in M. c = f x by auto
hence AE x∈A in M. False using assms(2) by auto
hence A ∈ null-sets M using AE-iff-null-sets A-in-sets-M by metis
thus False using A(3) by (simp add: emeasure-restr-to-subalg null-setsD1
subalg)
qed
thus ?thesis using AE-iff-null-sets[OF X-in-M] unfolding X-def by auto
qed

corollary cond-exp-less-c:
fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, or-
dered-real-vector}
assumes integrable M f AE x in M. f x < c
shows AE x in M. cond-exp M F f x < c
proof -
have AE x in M. cond-exp M F f x = - cond-exp M F (λx. - f x) x using
cond-exp-uminus[OF assms(1)] by auto
moreover have AE x in M. cond-exp M F (λx. - f x) x > - c using assms
by (intro cond-exp-gr-c) auto

```

**ultimately show** ?thesis **by** (force simp add: minus-less-iff)  
**qed**

**lemma** cond-exp-mono-strict:

**fixes**  $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$   
  **assumes** integrable  $M f$  integrable  $M g$  AE  $x$  in  $M$ .  $f x < g x$   
  **shows** AE  $x$  in  $M$ . cond-exp  $M F f x < cond-exp M F g x$   
  **using** cond-exp-less-c[*OF Bochner-Integration.integrable-diff*, *OF assms(1,2)*, of 0]  
    cond-exp-diff[*OF assms(1,2)*] *assms(3)* **by** auto

**lemma** cond-exp-ge-c:

**fixes**  $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$   
  **assumes** [measurable]: integrable  $M f$   
  **and** AE  $x$  in  $M$ .  $f x \geq c$   
  **shows** AE  $x$  in  $M$ . cond-exp  $M F f x \geq c$

**proof** –

  let  $?F = restr-to-subalg M F$   
  **interpret** sigma-finite-measure *restr-to-subalg*  $M F$  **using** sigma-fin-subalg **by** auto  
  {  
    fix  $A$  **assume**  $asm: A \in sets ?F$   $0 < measure ?F A$   
    **have** [simp]:  $sets ?F = sets F$   $measure ?F A = measure M A$  **using** *asm* **by** (auto simp add: measure-def sets-restr-to-subalg[*OF subalg*] emeasure-restr-to-subalg[*OF subalg*])  
    **have**  $M\text{-}A: emeasure M A < \infty$  **using** measure-zero-top *asm* **by** (force simp add: top.not-eq-extremum)  
    **hence**  $F\text{-}A: emeasure ?F A < \infty$  **using** *asm(1)* emeasure-restr-to-subalg *subalg* **by** fastforce  
    **have** set-lebesgue-integral  $M A (\lambda\_. c) \leq set-lebesgue-integral M A f$  **using** assms *asm M-A subalg* **by** (intro set-integral-mono-AE-banach, auto simp add: set-integrable-def integrable-mult-indicator subalgebra-def sets-restr-to-subalg)  
    **also have** ... = set-lebesgue-integral  $M A$  (cond-exp  $M F f$ ) **using** cond-exp-set-integral[*OF assms(1)*] *asm* **by** auto  
    **also have** ... = set-lebesgue-integral  $?F A$  (cond-exp  $M F f$ ) **unfolding** set-lebesgue-integral-def **using** *asm borel-measurable-cond-exp* **by** (intro integral-subalgebra2[*OF subalg, symmetric*], simp)  
    **finally have**  $(1 / measure ?F A) *_R set-lebesgue-integral ?F A (cond-exp M F f) \in \{c..\}$  **using** *asm subalg M-A* **by** (auto simp add: set-integral-const subalgebra-def intro!: pos-divideR-le-eq[THEN iffD1])  
  }  
  **thus** ?thesis **using** AE-restr-to-subalg[*OF subalg*] averaging-theorem[*OF integrable-in-subalg closed-atLeast, OF subalg borel-measurable-cond-exp integrable-cond-exp*] **by** auto  
**qed**

**corollary** cond-exp-le-c:

**fixes**  $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$   
**assumes** integrable  $M f$   
**and**  $\text{AE } x \text{ in } M. f x \leq c$   
**shows**  $\text{AE } x \text{ in } M. \text{cond-exp } M F f x \leq c$   
**proof –**  
**have**  $\text{AE } x \text{ in } M. \text{cond-exp } M F f x = - \text{cond-exp } M F (\lambda x. - f x) x$  **using**  
 $\text{cond-exp-uminus}[\text{OF assms}(1)]$  **by** force  
**moreover have**  $\text{AE } x \text{ in } M. \text{cond-exp } M F (\lambda x. - f x) x \geq -c$  **using** assms  
**by** (intro cond-exp-ge-c) auto  
**ultimately show** ?thesis **by** (force simp add: minus-le-iff)  
**qed**

**corollary** cond-exp-mono:

**fixes**  $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$   
**assumes** integrable  $M f$  integrable  $M g$   $\text{AE } x \text{ in } M. f x \leq g x$   
**shows**  $\text{AE } x \text{ in } M. \text{cond-exp } M F f x \leq \text{cond-exp } M F g x$   
**using** cond-exp-le-c[ $\text{OF Bochner-Integration.integrable-diff}$ , OF assms(1,2), of 0]  
 $\text{cond-exp-diff}[\text{OF assms}(1,2)]$  assms(3) **by** auto

**corollary** cond-exp-min:

**fixes**  $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$   
**assumes** integrable  $M f$  integrable  $M g$   
**shows**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min(f x) (g x)) \xi \leq \min(\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$   
**proof –**  
**have**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min(f x) (g x)) \xi \leq \text{cond-exp } M F f \xi$  **by**  
 (intro cond-exp-mono integrable-min assms, simp)  
**moreover have**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min(f x) (g x)) \xi \leq \text{cond-exp } M F g \xi$  **by** (intro cond-exp-mono integrable-min assms, simp)  
**ultimately show**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \min(f x) (g x)) \xi \leq \min(\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$  **by** fastforce  
**qed**

**corollary** cond-exp-max:

**fixes**  $f :: 'a \Rightarrow 'b :: \{second-countable-topology, banach, linorder-topology, ordered-real-vector\}$   
**assumes** integrable  $M f$  integrable  $M g$   
**shows**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max(f x) (g x)) \xi \geq \max(\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$   
**proof –**  
**have**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max(f x) (g x)) \xi \geq \text{cond-exp } M F f \xi$  **by**  
 (intro cond-exp-mono integrable-max assms, simp)  
**moreover have**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max(f x) (g x)) \xi \geq \text{cond-exp } M F g \xi$  **by** (intro cond-exp-mono integrable-max assms, simp)  
**ultimately show**  $\text{AE } \xi \text{ in } M. \text{cond-exp } M F (\lambda x. \max(f x) (g x)) \xi \geq \max(\text{cond-exp } M F f \xi) (\text{cond-exp } M F g \xi)$  **by** fastforce

```

(cond-exp M F f ξ) (cond-exp M F g ξ) by fastforce
qed

corollary cond-exp-inf:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}
  assumes integrable M f integrable M g
  shows AE ξ in M. cond-exp M F (λx. inf (f x) (g x)) ξ ≤ inf (cond-exp M F f
ξ) (cond-exp M F g ξ)
  unfolding inf-min using assms by (rule cond-exp-min)

corollary cond-exp-sup:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, linorder-topology, ordered-real-vector, lattice}
  assumes integrable M f integrable M g
  shows AE ξ in M. cond-exp M F (λx. sup (f x) (g x)) ξ ≥ sup (cond-exp M F f
ξ) (cond-exp M F g ξ)
  unfolding sup-max using assms by (rule cond-exp-max)

end

```

## 2.4 Probability Spaces

```

lemma (in prob-space) sigma-finite-subalgebra-restr-to-subalg:
  assumes subalgebra M F
  shows sigma-finite-subalgebra M F
proof (intro sigma-finite-subalgebra.intro)
  interpret F: prob-space restr-to-subalg M F using assms prob-space-restr-to-subalg
  prob-space-axioms by blast
  show sigma-finite-measure (restr-to-subalg M F) by (rule F.sigma-finite-measure-axioms)
qed (rule assms)

lemma (in prob-space) cond-exp-trivial:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach}
  assumes integrable M f
  shows AE x in M. cond-exp M (sigma (space M) {}) f x = expectation f
proof -
  interpret sigma-finite-subalgebra M sigma (space M) {} by (auto intro: sigma-finite-subalgebra-restr-to-subalg
  simp add: subalgebra-def sigma-sets-empty-eq)
  show ?thesis using assms by (intro cond-exp-charact) (auto simp add: sigma-sets-empty-eq
  set-lebesgue-integral-def prob-space cong: Bochner-Integration.integral-cong)
qed

```

The following lemma shows that independent  $\sigma$ -algebras don't matter for the conditional expectation. The proof is adapted from [4].

```

lemma (in prob-space) cond-exp-indep-subalgebra:
  fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, real-normed-field}
  assumes subalgebra: subalgebra M F subalgebra M G
  and independent: indep-set G (sigma (space M) (F ∪ vimage-algebra (space

```

```

 $M) f borel))$ 
assumes [measurable]: integrable  $M f$ 
shows  $\text{AE } x \text{ in } M. \text{ cond-exp } M (\text{sigma (space } M) (F \cup G)) f x = \text{cond-exp } M F f x$ 
proof -
  interpret  $\text{Un-sigma: sigma-finite-subalgebra } M \text{ sigma (space } M) (F \cup G)$  using
   $\text{assms}(1,2)$  by (auto intro!: sigma-finite-subalgebra-restr-to-subalg sets.sigma-sets-subset
  simp add: subalgebra-def space-measure-of-conv sets-measure-of-conv)
  interpret  $\text{sigma-finite-subalgebra } M F$  using  $\text{assms}$  by (auto intro: sigma-finite-subalgebra-restr-to-subalg)
  {
    fix  $A$ 
    assume  $\text{asm: } A \in \text{sigma (space } M) \{a \cap b \mid a b. a \in F \wedge b \in G\}$ 
    have  $\text{in-events: sigma-sets (space } M) \{a \cap b \mid a b. a \in \text{sets } F \wedge b \in \text{sets } G\} \subseteq \text{events}$  using subalgebra by (intro sets.sigma-sets-subset, auto simp add:
    subalgebra-def)
    have  $\text{Int-stable } \{a \cap b \mid a b. a \in F \wedge b \in G\}$ 
    proof -
    {
      fix  $af bf ag bg$ 
      assume  $F: af \in F bf \in F$  and  $G: ag \in G bg \in G$ 
      have  $af \cap bf \in F$  by (intro sets.Int F)
      moreover have  $ag \cap bg \in G$  by (intro sets.Int G)
      ultimately have  $\exists a b. af \cap ag \cap (bf \cap bg) = a \cap b \wedge a \in \text{sets } F \wedge b \in \text{sets } G$  by (metis inf-assoc inf-left-commute)
    }
    thus ?thesis by (force intro!: Int-stableI)
  qed
  moreover have  $\{a \cap b \mid a b. a \in F \wedge b \in G\} \subseteq \text{Pow (space } M)$  using
  subalgebra by (force simp add: subalgebra-def dest: sets.sets-into-space)
  moreover have  $A \in \text{sigma-sets (space } M) \{a \cap b \mid a b. a \in F \wedge b \in G\}$  using
  calculation  $\text{asm by force}$ 
  ultimately have  $\text{set-lebesgue-integral } M A f = \text{set-lebesgue-integral } M A$ 
  ( $\text{cond-exp } M F f$ )
  proof (induction rule: sigma-sets-induct-disjoint)
  case (basic  $A$ )
  then obtain  $a b$  where  $A: A = a \cap b a \in F b \in G$  by blast
  hence events[measurable]:  $a \in \text{events } b \in \text{events}$  using subalgebra by (auto
  simp add: subalgebra-def)
  have [simp]:  $\text{sigma-sets (space } M) \{\text{indicator } b -` A \cap \text{space } M \mid A. A \in \text{borel}\}$ 
   $\subseteq G$ 
  using borel-measurable-indicator[OF A(3), THEN measurable-sets] sets.top
  subalgebra
  by (intro sets.sigma-sets-subset') (fastforce simp add: subalgebra-def) +
  have  $\text{Un-in-sigma: } F \cup \text{vimage-algebra (space } M) f borel \subseteq \text{sigma (space } M) (F$ 
   $\cup \text{vimage-algebra (space } M) f borel)$  by (metis equalityE le-supI sets.space-closed
  sigma-le-sets space-vimage-algebra subalg subalgebra-def)

```

```

have [intro]: indep-var borel (indicator b) borel ( $\lambda\omega.$  indicator a  $\omega *_R f \omega$ )
proof -
  have [simp]: sigma-sets (space M)  $\{(\lambda\omega.$  indicator a  $\omega *_R f \omega) -' A \cap space$ 
 $M |A. A \in borel\} \subseteq sigma (space M)$  ( $F \cup vimage-algebra (space M) f borel$ )
  proof -
    have *:  $(\lambda\omega.$  indicator a  $\omega *_R f \omega) \in borel-measurable (sigma (space M)$ 
 $(F \cup vimage-algebra (space M) f borel))$ 
    using borel-measurable-indicator[OF A(2), THEN measurable-sets, OF
borel-open] subalgebra
    by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN
subsetD])
      (auto simp add: space-measure-of-conv subalgebra-def sets-vimage-algebra2)
      thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',
auto simp add: space-measure-of-conv)
    qed
    have indep-set (sigma-sets (space M) {indicator b -' A \cap space M |A. A \in borel}) ( $\lambda\omega.$  indicator a  $\omega *_R f \omega) -' A \cap space M |A. A \in borel\}$ 
    using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
    thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)
    qed

    have [intro]: indep-var borel (indicator b) borel ( $\lambda\omega.$  indicat-real a  $\omega *_R$ 
cond-exp M F f \omega)
    proof -
      have [simp]: sigma-sets (space M)  $\{(\lambda\omega.$  indicator a  $\omega *_R cond-exp M F f$ 
 $\omega) -' A \cap space M |A. A \in borel\} \subseteq sigma (space M)$  ( $F \cup vimage-algebra (space$ 
 $M) f borel$ )
      proof -
        have *:  $(\lambda\omega.$  indicator a  $\omega *_R cond-exp M F f \omega) \in borel-measurable (sigma$ 
 $(space M) (F \cup vimage-algebra (space M) f borel))$ 
        using borel-measurable-indicator[OF A(2), THEN measurable-sets, OF
borel-open] subalgebra
          borel-measurable-cond-exp[THEN measurable-sets, OF borel-open, of
 $- M F f]$ 
          by (intro borel-measurable-scaleR borel-measurableI Un-in-sigma[THEN
subsetD])
            (auto simp add: space-measure-of-conv subalgebra-def)
            thus ?thesis using measurable-sets[OF *] by (intro sets.sigma-sets-subset',
auto simp add: space-measure-of-conv)
        qed
        have indep-set (sigma-sets (space M) {indicator b -' A \cap space M |A. A \in borel}) ( $\lambda\omega.$  indicator a  $\omega *_R cond-exp M F f \omega) -' A \cap$ 
space M |A. A \in borel\}
        using independent unfolding indep-set-def by (rule indep-sets-mono-sets,
auto split: bool.split)
        thus ?thesis by (subst indep-var-eq, auto intro!: borel-measurable-scaleR)

```

```

qed

have set-lebesgue-integral M A f = (LINT x|M. indicator b x * (indicator a
x *_R f x))
  unfolding set-lebesgue-integral-def A indicator-inter-arith
  by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
  also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *_R f
x)
  by (intro indep-var-lebesgue-integral
        Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
        integrable-mult-indicator[OF - assms(4)], blast) (auto simp add:
indicator-def)
  also have ... = (LINT x|M. indicator b x) * (LINT x|M. indicator a x *_R
cond-exp M F f x)
  using cond-exp-set-integral[OF assms(4) A(2)] unfolding set-lebesgue-integral-def
by argo
  also have ... = (LINT x|M. indicator b x * (indicator a x *_R cond-exp M
F f x))
  by (intro indep-var-lebesgue-integral[symmetric]
        Bochner-Integration.integrable-bound[OF integrable-const[of 1 :: 'b]
borel-measurable-indicator]
        integrable-mult-indicator[OF - integrable-cond-exp], blast) (auto simp
add: indicator-def)
  also have ... = set-lebesgue-integral M A (cond-exp M F f)
  unfolding set-lebesgue-integral-def A indicator-inter-arith
  by (intro Bochner-Integration.integral-cong, auto simp add: scaleR-scaleR[symmetric]
indicator-times-eq-if(1))
  finally show ?case .
next
  case empty
  then show ?case unfolding set-lebesgue-integral-def by simp
next
  case (compl A)
  have A-in-space: A ⊆ space M using compl using in-events sets.sets-into-space
  by blast
  have set-lebesgue-integral M (space M - A) f = set-lebesgue-integral M (space
M - A ∪ A) f - set-lebesgue-integral M A f
    using compl(1) in-events
    by (subst set-integral-Un[of space M - A A], blast)
      (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
assms(4)], fast)+
  also have ... = set-lebesgue-integral M (space M - A ∪ A) (cond-exp M F f)
  - set-lebesgue-integral M A (cond-exp M F f)
    using cond-exp-set-integral[OF assms(4) sets.top] compl subalgebra by (simp
add: subalgebra-def Un-absorb2[OF A-in-space])
  also have ... = set-lebesgue-integral M (space M - A) (cond-exp M F f)
    using compl(1) in-events

```

```

by (subst set-integral-Un[of space M - A A], blast)
  (simp | intro integrable-mult-indicator[folded set-integrable-def, OF -
integrable-cond-exp], fast) +
  finally show ?case .
next
  case (union A)
  have set-lebesgue-integral M (UN (range A)) f = (∑ i. set-lebesgue-integral M
(A i) f)
    using union in-events
    by (intro lebesgue-integral-countable-add) (auto simp add: disjoint-family-onD
intro!: integrable-mult-indicator[folded set-integrable-def, OF - assms(4)])
    also have ... = (∑ i. set-lebesgue-integral M (A i) (cond-exp M F f)) using
union by presburger
    also have ... = set-lebesgue-integral M (UN (range A)) (cond-exp M F f)
    using union in-events
    by (intro lebesgue-integral-countable-add[symmetric]) (auto simp add: dis-
joint-family-onD intro!: integrable-mult-indicator[folded set-integrable-def, OF - in-
tegrable-cond-exp])
    finally show ?case .
qed
}
moreover have sigma (space M) {a ∩ b | a b. a ∈ F ∧ b ∈ G} = sigma (space
M) (F ∪ G)
proof -
  have sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G} = sigma-sets
(space M) (sets F ∪ sets G)
  proof -
    fix a b assume asm: a ∈ F b ∈ G
    hence a ∩ b ∈ sigma-sets (space M) (F ∪ G) using subalgebra unfolding
Int-range-binary by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) (force simp
add: subalgebra-def dest: sets.sets-into-space) +
  }
  moreover
  {
    fix a
    assume a ∈ sets F
    hence a ∈ sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G}
      using subalgebra sets.top[of G] sets.sets-into-space[of - F]
      by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
  }
  moreover
  {
    fix a assume a ∈ sets F ∨ a ∈ sets G a ∉ sets F
    hence a ∈ sets G by blast
    hence a ∈ sigma-sets (space M) {a ∩ b | a b. a ∈ sets F ∧ b ∈ sets G}
      using subalgebra sets.top[of F] sets.sets-into-space[of - G]
      by (intro sigma-sets.Basic, auto simp add: subalgebra-def)
  }
}

```

```

ultimately show ?thesis by (intro sigma-sets-eqI) auto
qed
thus ?thesis using subalgebra by (intro sigma-eqI) (force simp add: subalgebra-def dest: sets.sets-into-space)+  

qed
moreover have (cond-exp M F f) ∈ borel-measurable (sigma (space M) (sets F  

∪ sets G))
proof -
have F ⊆ sigma (space M) (F ∪ G) by (metis Un-least Un-upper1  

measure-of-measure sets.space-closed sets-measure-of sigma-sets-subseteq subalg sub-  

algebra(2) subalgebra-def)
thus ?thesis using borel-measurable-cond-exp[THEN measurable-sets, OF borel-open,  

of - M F f] subalgebra by (intro borel-measurableI, force simp only: space-measure-of-conv  

subalgebra-def)
qed
ultimately show ?thesis using assms(4) integrable-cond-exp by (intro Un-sigma.cond-exp-charact)  

presburger+
qed

```

If a random variable is independent of a  $\sigma$ -algebra  $F$ , its conditional expectation  $\text{cond-exp } M F f$  is just its expectation.

```

lemma (in prob-space) cond-exp-indep:
fixes f :: 'a ⇒ 'b :: {second-countable-topology, banach, real-normed-field}
assumes subalgebra: subalgebra M F
and independent: indep-set F (vimage-algebra (space M) f borel)
and integrable: integrable M f
shows AE x in M. cond-exp M F f x = expectation f
proof -
have indep-set F (sigma (space M) (sigma (space M) {} ∪ (vimage-algebra (space  

M) f borel)))
using independent unfolding indep-set-def
by (rule indep-sets-mono-sets, simp add: bool.split)
(metis bot.extremum dual-order.refl sets.sets-measure-of-eq sets.sigma-sets-subset'  

sets-vimage-algebra-space space-vimage-algebra sup.absorb-iff2)
hence cond-exp-indep: AE x in M. cond-exp M (sigma (space M) (sigma (space  

M) {} ∪ F)) f x = expectation f
using cond-exp-indep-subalgebra[OF - subalgebra - integrable, of sigma (space  

M) {}] cond-exp-trivial[OF integrable]
by (auto simp add: subalgebra-def sigma-sets-empty-eq)
have sets (sigma (space M) (sigma (space M) {} ∪ F)) = F
using subalgebra sets.top[of F] unfolding subalgebra-def
by (simp add: sigma-sets-empty-eq, subst insert-absorb[of space M F], blast)
(metis insert-absorb[OF sets.empty-sets] sets.sets-measure-of-eq)
hence AE x in M. cond-exp M (sigma (space M) (sigma (space M) {} ∪ F)) f  

x = cond-exp M F f x by (rule cond-exp-sets-cong)
thus ?thesis using cond-exp-indep by force
qed

```

end

```

theory Filtered-Measure
imports HOL-Probability.Conditional-Expectation
begin

3 Filtered Measure Spaces

3.1 Filtered Measure

locale filtered-measure =
fixes M F and t0 :: 'b :: {second-countable-topology, order-topology, t2-space}
assumes subalgebras:  $\bigwedge i. t_0 \leq i \implies \text{subalgebra } M (F i)$ 
and sets-F-mono:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{sets } (F i) \leq \text{sets } (F j)$ 
begin

lemma space-F[simp]:
assumes t0 ≤ i
shows space (F i) = space M
using subalgebras assms by (simp add: subalgebra-def)

lemma sets-F-subset[simp]:
assumes t0 ≤ i
shows sets (F i) ⊆ sets M
using subalgebras assms by (simp add: subalgebra-def)

lemma subalgebra-F[intro]:
assumes t0 ≤ i i ≤ j
shows subalgebra (F j) (F i)
unfolding subalgebra-def using assms by (simp add: sets-F-mono)

lemma borel-measurable-mono:
assumes t0 ≤ i i ≤ j
shows borel-measurable (F i) ⊆ borel-measurable (F j)
unfolding subset-iff by (metis assms subalgebra-F measurable-from-subalg)

end

locale linearly-filtered-measure = filtered-measure M F t0 for M and F :: - :: {linorder-topology, conditionally-complete-lattice} ⇒ - and t0

context linearly-filtered-measure
begin

σ-algebra at infinity

definition F-infinity :: 'a measure where
F-infinity = sigma (space M) (UNION t ∈ {t0..}. sets (F t))

notation F-infinity ((F∞))

```

```

lemma space-F-infinity[simp]: space F∞ = space M unfolding F-infinity-def space-measure-of-conv
by simp

lemma sets-F-infinity: sets F∞ = sigma-sets (space M) ( $\bigcup t \in \{t_0..\}. \text{sets } (F t)$ )
  unfolding F-infinity-def using sets.space-closed[of F -] space-F by (blast intro!
  sets-measure-of)

lemma subset-F-infinity:
  assumes t ≥ t0
  shows F t ⊆ F∞ unfolding sets-F-infinity using assms by blast

lemma F-infinity-subset: F∞ ⊆ M
  unfolding sets-F-infinity using sets-F-subset
  by (simp add: SUP-le-iff sets.sigma-sets-subset)

lemma F-infinity-measurableI:
  assumes t ≥ t0 f ∈ borel-measurable (F t)
  shows f ∈ borel-measurable (F∞)
  by (metis assms borel-measurable-subalgebra space-F space-F-infinity subset-F-infinity)

end

locale nat-filtered-measure = linearly-filtered-measure M F 0 for M and F :: nat
  ⇒ -
locale enat-filtered-measure = linearly-filtered-measure M F 0 for M and F :: enat
  ⇒ -
locale real-filtered-measure = linearly-filtered-measure M F 0 for M and F :: real
  ⇒ -
locale ennreal-filtered-measure = linearly-filtered-measure M F 0 for M and F :: ennreal
  ⇒ -

```

### 3.2 σ-Finite Filtered Measure

The locale presented here is a generalization of the *sigma-finite-subalgebra* for a particular filtration.

```

locale sigma-finite-filtered-measure = filtered-measure +
  assumes sigma-finite-initial: sigma-finite-subalgebra M (F t0)

lemma (in sigma-finite-filtered-measure) sigma-finite-subalgebra-F[intro]:
  assumes t0 ≤ i
  shows sigma-finite-subalgebra M (F i)
  using assms by (metis dual-order.refl sets-F-mono sigma-finite-initial sigma-finite-subalgebra.nested-subalg-is-
  subalgebras subalgebra-def)

locale nat-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 :: nat for M F
locale enat-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 :: enat for M F

```

```

locale real-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0 ::  

real for M F
locale ennreal-sigma-finite-filtered-measure = sigma-finite-filtered-measure M F 0  

:: ennreal for M F

sublocale nat-sigma-finite-filtered-measure ⊆ nat-filtered-measure ..
sublocale enat-sigma-finite-filtered-measure ⊆ enat-filtered-measure ..
sublocale real-sigma-finite-filtered-measure ⊆ real-filtered-measure ..
sublocale ennreal-sigma-finite-filtered-measure ⊆ ennreal-filtered-measure ..

sublocale nat-sigma-finite-filtered-measure ⊆ sigma-finite-subalgebra M F i by  

blast
sublocale enat-sigma-finite-filtered-measure ⊆ sigma-finite-subalgebra M F i by  

fastforce
sublocale real-sigma-finite-filtered-measure ⊆ sigma-finite-subalgebra M F |i| by  

fastforce
sublocale ennreal-sigma-finite-filtered-measure ⊆ sigma-finite-subalgebra M F i by  

fastforce

```

### 3.3 Finite Filtered Measure

```

locale finite-filtered-measure = filtered-measure + finite-measure

sublocale finite-filtered-measure ⊆ sigma-finite-filtered-measure
using subalgebras by (unfold-locales, blast, meson dual-order.refl finite-measure-axioms
finite-measure-def finite-measure-restr-to-subalg sigma-finite-measure.sigma-finite-countable)

locale nat-finite-filtered-measure = finite-filtered-measure M F 0 :: nat for M F
locale enat-finite-filtered-measure = finite-filtered-measure M F 0 :: enat for M F
locale real-finite-filtered-measure = finite-filtered-measure M F 0 :: real for M F
locale ennreal-finite-filtered-measure = finite-filtered-measure M F 0 :: ennreal for
M F

sublocale nat-finite-filtered-measure ⊆ nat-sigma-finite-filtered-measure ..
sublocale enat-finite-filtered-measure ⊆ enat-sigma-finite-filtered-measure ..
sublocale real-finite-filtered-measure ⊆ real-sigma-finite-filtered-measure ..
sublocale ennreal-finite-filtered-measure ⊆ ennreal-sigma-finite-filtered-measure ..

```

### 3.4 Constant Filtration

```

lemma filtered-measure-constant-filtration:
assumes subalgebra M F
shows filtered-measure M (λ-. F) t₀
using assms by (unfold-locales) blast+

sublocale sigma-finite-subalgebra ⊆ constant-filtration: sigma-finite-filtered-measure
M λ- :: 't :: {second-countable-topology, linorder-topology}. F t₀
using subalg by (unfold-locales) blast+

lemma (in finite-measure) filtered-measure-constant-filtration:

```

```

assumes subalgebra M F
shows finite-filtered-measure M (λ-. F) t₀
using assms by (unfold-locales) blast+
end

```

```

theory Stochastic-Process
imports Filtered-Measure Measure-Space-Supplement HOL-Probability.Independent-Family
begin

```

## 4 Stochastic Processes

### 4.1 Stochastic Process

A stochastic process is a collection of random variables, indexed by a type ' $b$ '.

```

locale stochastic-process =
fixes M t₀ and X :: 'b :: {second-countable-topology, order-topology, t2-space} ⇒
'a ⇒ 'c :: {second-countable-topology, banach}
assumes random-variable[measurable]: ∀ i. t₀ ≤ i ⇒ X i ∈ borel-measurable M
begin

definition left-continuous where left-continuous = (AE ξ in M. ∀ t. continuous
(at-left t) (λ i. X i ξ))
definition right-continuous where right-continuous = (AE ξ in M. ∀ t. continuous
(at-right t) (λ i. X i ξ))

end

```

```

lemma stochastic-process-const-fun:
assumes f ∈ borel-measurable M
shows stochastic-process M t₀ (λ i. f) using assms by (unfold-locales)

```

```

lemma stochastic-process-const:
shows stochastic-process M t₀ (λ i. c i) by (unfold-locales) simp

```

In the following segment, we cover basic operations on stochastic processes.

```

context stochastic-process
begin

```

```

lemma compose-stochastic:
assumes ∀ i. t₀ ≤ i ⇒ f i ∈ borel-measurable borel
shows stochastic-process M t₀ (λ i. ξ. (f i) (X i ξ))
by (unfold-locales) (intro measurable-compose[OF random-variable assms])

```

```

lemma norm-stochastic: stochastic-process M t₀ (λ i. ξ. norm (X i ξ)) by (fastforce
intro: compose-stochastic)

```

```

lemma scaleR-right-stochastic:
  assumes stochastic-process M t0 Y
  shows stochastic-process M t0 ( $\lambda i \xi. (Y i \xi) *_R (X i \xi)$ )
  using stochastic-process.random-variable[OF assms] random-variable by (unfold-locales)
  simp

lemma scaleR-right-const-fun-stochastic:
  assumes f ∈ borel-measurable M
  shows stochastic-process M t0 ( $\lambda i \xi. f \xi *_R (X i \xi)$ )
  by (unfold-locales) (intro borel-measurable-scaleR assms random-variable)

lemma scaleR-right-const-stochastic: stochastic-process M t0 ( $\lambda i \xi. c i *_R (X i \xi)$ )
  by (unfold-locales) simp

lemma add-stochastic:
  assumes stochastic-process M t0 Y
  shows stochastic-process M t0 ( $\lambda i \xi. X i \xi + Y i \xi$ )
  using stochastic-process.random-variable[OF assms] random-variable by (unfold-locales)
  simp

lemma diff-stochastic:
  assumes stochastic-process M t0 Y
  shows stochastic-process M t0 ( $\lambda i \xi. X i \xi - Y i \xi$ )
  using stochastic-process.random-variable[OF assms] random-variable by (unfold-locales)
  simp

lemma uminus-stochastic: stochastic-process M t0 (-X) using scaleR-right-const-stochastic[of λ-. -1] by (simp add: fun-Compl-def)

lemma partial-sum-stochastic: stochastic-process M t0 ( $\lambda n \xi. \sum_{i \in \{t_0..n\}} X i \xi$ )
  by (unfold-locales) simp

lemma partial-sum'-stochastic: stochastic-process M t0 ( $\lambda n \xi. \sum_{i \in \{t_0..<n\}} X i \xi$ )
  by (unfold-locales) simp

end

lemma stochastic-process-sum:
  assumes  $\bigwedge i. i \in I \implies$  stochastic-process M t0 (X i)
  shows stochastic-process M t0 ( $\lambda k \xi. \sum_{i \in I} X i k \xi$ ) using assms[THEN stochastic-process.random-variable]] by (unfold-locales, auto)

```

**4.1.1 Natural Filtration**

The natural filtration induced by a stochastic process  $X$  is the filtration generated by all events involving the process up to the time index  $t$ , i.e.  $F_t = \sigma(\{X s \mid s. s \leq t\})$ .

**definition** natural-filtration :: '*a* measure ⇒ '*b* ⇒ ('*b* ⇒ '*a* ⇒ '*c* :: topologi-

*cal-space)  $\Rightarrow$  'b :: {second-countable-topology, order-topology}  $\Rightarrow$  'a measure where  
natural-filtration M t<sub>0</sub> Y = ( $\lambda t$ . family-vimage-algebra (space M) {Y i | i. i  $\in$  {t<sub>0</sub>..t}}) borel)*

**abbreviation** nat-natural-filtration  $\equiv \lambda M$ . natural-filtration M (0 :: nat)  
**abbreviation** real-natural-filtration  $\equiv \lambda M$ . natural-filtration M (0 :: real)

**lemma** space-natural-filtration[simp]: space (natural-filtration M t<sub>0</sub> X t) = space M **unfolding** natural-filtration-def space-family-vimage-algebra ..

**lemma** sets-natural-filtration: sets (natural-filtration M t<sub>0</sub> X t) = sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}}$ . {X i - ' A  $\cap$  space M | A. A  $\in$  borel})  
**unfolding** natural-filtration-def sets-family-vimage-algebra **by** (intro sigma-sets-eqI)  
blast+

**lemma** sets-natural-filtration':  
**assumes** borel = sigma UNIV S  
**shows** sets (natural-filtration M t<sub>0</sub> X t) = sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}}$ . {X i - ' A  $\cap$  space M | A. A  $\in$  S})  
**proof** (subst sets-natural-filtration, intro sigma-sets-eqI, clarify)  
fix i and A :: 'a set **assume** asm: i  $\in$  {t<sub>0</sub>..t} A  $\in$  sets borel  
hence A  $\in$  sigma-sets UNIV S **unfolding** assms **by** simp  
thus X i - ' A  $\cap$  space M  $\in$  sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}}$ . {X i - ' A  $\cap$  space M | A. A  $\in$  S})  
**proof** (induction)  
**case** (Compl a)  
have X i - ' (UNIV - a)  $\cap$  space M = space M - (X i - ' a  $\cap$  space M) **by** blast  
**then show** ?case **using** Compl(2)[THEN sigma-sets.Compl] **by** presburger  
**next**  
**case** (Union a)  
have X i - '  $\bigcup$  (range a)  $\cap$  space M =  $\bigcup$  (range ( $\lambda j$ . X i - ' a j  $\cap$  space M))  
**by** blast  
**then show** ?case **using** Union(2)[THEN sigma-sets.Union] **by** presburger  
**qed** (auto intro: asm sigma-sets.Empty)  
**qed** (intro sigma-sets.Basic, force simp add: assms)

**lemma** sets-natural-filtration-open:  
sets (natural-filtration M t<sub>0</sub> X t) = sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}}$ . {X i - ' A  $\cap$  space M | A. open A})  
**using** sets-natural-filtration' **by** (force simp only: borel-def mem-Collect-eq)

**lemma** sets-natural-filtration-oi:  
sets (natural-filtration M t<sub>0</sub> X t) = sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}}$ . {X i - ' A  $\cap$  space M | A :: - :: {linorder-topology, second-countable-topology} set. A  $\in$  range greaterThan})  
**by** (rule sets-natural-filtration'[OF borel-Ioi])

**lemma** sets-natural-filtration-io:

*sets (natural-filtration M t<sub>0</sub> X t) = sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}} \{X i -` A \cap space M | A :: - :: \{linorder-topology, second-countable-topology\} set. A \in range lessThan\}$ )*

**by** (rule sets-natural-filtration'[OF borel-Iio])

**lemma** sets-natural-filtration-ci:

*sets (natural-filtration M t<sub>0</sub> X t) = sigma-sets (space M) ( $\bigcup_{i \in \{t_0..t\}} \{X i -` A \cap space M | A :: real set. A \in range atLeast\}$ )*

**by** (rule sets-natural-filtration'[OF borel-Ici])

**context** stochastic-process

**begin**

**lemma** subalgebra-natural-filtration:

**shows** subalgebra M (natural-filtration M t<sub>0</sub> X i)

**unfolding** subalgebra-def **using** measurable-family-iff-sets **by** (force simp add: natural-filtration-def)

**lemma** filtered-measure-natural-filtration:

**shows** filtered-measure M (natural-filtration M t<sub>0</sub> X) t<sub>0</sub>

**by** (unfold-locales) (intro subalgebra-natural-filtration, simp only: sets-natural-filtration, intro sigma-sets-subseteq, force)

In order to show that the natural filtration constitutes a filtered  $\sigma$ -finite measure, we need to provide a countable exhausting set in the preimage of  $X t_0$ .

**lemma** sigma-finite-filtered-measure-natural-filtration:

**assumes** exhausting-set: countable A ( $\bigcup A = space M \wedge \forall a \in A \implies emeasure M a \neq \infty \wedge \forall a \in A \implies \exists b \in borel. a = X t_0 -` b \cap space M$ )

**shows** sigma-finite-filtered-measure M (natural-filtration M t<sub>0</sub> X) t<sub>0</sub>

**proof** (unfold-locales)

**have** A  $\subseteq$  sets (restr-to-subalg M (natural-filtration M t<sub>0</sub> X t<sub>0</sub>)) **using** exhausting-set **by** (simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] sets-natural-filtration) fast

**moreover have**  $\bigcup A = space (restr-to-subalg M (natural-filtration M t_0 X t_0))$  **unfolding** space-restr-to-subalg **using** exhausting-set **by** simp

**moreover have**  $\forall a \in A. emeasure (restr-to-subalg M (natural-filtration M t_0 X t_0)) a \neq \infty$  **using** calculation(1) exhausting-set(3)

**by** (auto simp add: sets-restr-to-subalg[OF subalgebra-natural-filtration] emeasure-restr-to-subalg[OF subalgebra-natural-filtration])

**ultimately show**  $\exists A. countable A \wedge A \subseteq sets (restr-to-subalg M (natural-filtration M t_0 X t_0)) \wedge \bigcup A = space (restr-to-subalg M (natural-filtration M t_0 X t_0)) \wedge (\forall a \in A. emeasure (restr-to-subalg M (natural-filtration M t_0 X t_0)) a \neq \infty)$  **using** exhausting-set **by** blast

**show**  $\bigwedge i j. [t_0 \leq i; i \leq j] \implies sets (natural-filtration M t_0 X i) \subseteq sets (natural-filtration M t_0 X j)$  **using** filtered-measure.subalgebra-F[OF filtered-measure-natural-filtration]

**by** (simp add: subalgebra-def)

**qed** (auto intro: subalgebra-natural-filtration)

```

lemma finite-filtered-measure-natural-filtration:
  assumes finite-measure M
  shows finite-filtered-measure M (natural-filtration M t0 X) t0
  using finite-measure.axioms[OF assms] filtered-measure-natural-filtration by intro-locales

end

Filtration generated by independent variables.

lemma (in prob-space) indep-set-natural-filtration:
  assumes t0 ≤ s s < t indep-vars (λ-. borel) X {t0..}
  shows indep-set (natural-filtration M t0 X s) (vimage-algebra (space M) (X t)
borel)
proof -
  have indep-sets (λi. {X i -` A ∩ space M | A. A ∈ sets borel}) (UN (range (case-bool
{t0..s} {t})))
  using assms
  by (intro assms(3)[unfolded indep-vars-def, THEN conjunct2, THEN indep-sets-mono])
(auto simp add: case-bool-if)
  thus ?thesis unfolding indep-set-def using assms
  by (intro indep-sets-cong[THEN iffD1, OF refl -` indep-sets-collect-sigma[of λi.
{X i -` A ∩ space M | A. A ∈ borel} case-bool {t0..s} {t}]])
  (simp add: sets-natural-filtration sets-vimage-algebra split: bool.split, simp,
intro Int-stableI, clarsimp, metis sets.Int vimage-Int Int-commute Int-left-absorb
Int-left-commute, force simp add: disjoint-family-on-def split: bool.split)
qed

```

## 4.2 Adapted Process

We call a collection a stochastic process  $X$  adapted if  $X_i$  is  $F_i$ -borel-measurable for all indices  $i$ .

```

locale adapted-process = filtered-measure M F t0 for M F t0 and X :: - ⇒ - ⇒ -
:: {second-countable-topology, banach} +
assumes adapted[measurable]: ∀i. t0 ≤ i ⇒ X i ∈ borel-measurable (F i)
begin

lemma adaptedE[elim]:
  assumes [!∀j i. t0 ≤ j ⇒ j ≤ i ⇒ X j ∈ borel-measurable (F i)] ⇒ P
  shows P
  using assms using adapted by (metis dual-order.trans borel-measurable-subalgebra
sets-F-mono space-F)

lemma adaptedD:
  assumes t0 ≤ j j ≤ i
  shows X j ∈ borel-measurable (F i) using assms adaptedE by meson

end

```

```

lemma (in filtered-measure) adapted-process-const-fun:
  assumes  $f \in \text{borel-measurable } (F t_0)$ 
  shows adapted-process  $M F t_0 (\lambda i. f)$ 
  using measurable-from-subalg subalgebra- $F$  assms by (unfold-locales) blast

```

```

lemma (in filtered-measure) adapted-process-const:
  shows adapted-process  $M F t_0 (\lambda i. c i)$  by (unfold-locales) simp

```

Again, we cover basic operations.

```

context adapted-process
begin

```

```

lemma compose-adapted:
  assumes  $\bigwedge i. t_0 \leq i \implies f i \in \text{borel-measurable borel}$ 
  shows adapted-process  $M F t_0 (\lambda i. \xi. (f i) (X i \xi))$ 
  by (unfold-locales) (intro measurable-compose[OF adapted assms])

```

```

lemma norm-adapted: adapted-process  $M F t_0 (\lambda i. \xi. \text{norm} (X i \xi))$  by (fastforce
  intro: compose-adapted)

```

```

lemma scaleR-right-adapted:
  assumes adapted-process  $M F t_0 R$ 
  shows adapted-process  $M F t_0 (\lambda i. \xi. (R i \xi) *_R (X i \xi))$ 
  using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp

```

```

lemma scaleR-right-const-fun-adapted:
  assumes  $f \in \text{borel-measurable } (F t_0)$ 
  shows adapted-process  $M F t_0 (\lambda i. \xi. f \xi *_R (X i \xi))$ 
  using assms by (fast intro: scaleR-right-adapted adapted-process-const-fun)

```

```

lemma scaleR-right-const-adapted: adapted-process  $M F t_0 (\lambda i. \xi. c i *_R (X i \xi))$ 
  by (unfold-locales) simp

```

```

lemma add-adapted:
  assumes adapted-process  $M F t_0 Y$ 
  shows adapted-process  $M F t_0 (\lambda i. \xi. X i \xi + Y i \xi)$ 
  using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp

```

```

lemma diff-adapted:
  assumes adapted-process  $M F t_0 Y$ 
  shows adapted-process  $M F t_0 (\lambda i. \xi. X i \xi - Y i \xi)$ 
  using adapted-process.adapted[OF assms] adapted by (unfold-locales) simp

```

```

lemma uminus-adapted: adapted-process  $M F t_0 (-X)$  using scaleR-right-const-adapted[of
   $\lambda i. -1$ ] by (simp add: fun-Compl-def)

```

```

lemma partial-sum-adapted: adapted-process  $M F t_0 (\lambda n. \xi. \sum_{i \in \{t_0..n\}} X i \xi)$ 
proof (unfold-locales)
  fix  $i :: 'b$ 

```

```

have X j ∈ borel-measurable (F i) if  $t_0 \leq j \leq i$  for j using that adaptedE by
meson
  thus ( $\lambda\xi. \sum_{i \in \{t_0..i\}} X i \xi$ ) ∈ borel-measurable (F i) by simp
qed

lemma partial-sum'-adapted: adapted-process M F t₀ ( $\lambda n \xi. \sum_{i \in \{t_0..<n\}} X i \xi$ )

proof (unfold-locales)
  fix i :: 'b
  have X j ∈ borel-measurable (F i) if  $t_0 \leq j < i$  for j using that adaptedE by
fastforce
  thus ( $\lambda\xi. \sum_{i \in \{t_0..<i\}} X i \xi$ ) ∈ borel-measurable (F i) by simp
qed

end

```

In the discrete time case, we have the following lemmas which will be useful later on.

```

lemma (in nat-filtered-measure) partial-sum-Suc-adapted:
  assumes adapted-process M F 0 X
  shows adapted-process M F 0 ( $\lambda n \xi. \sum_{i < n} X (\text{Suc } i) \xi$ )
proof (unfold-locales)
  interpret adapted-process M F 0 X using assms by blast
  fix i
  have X j ∈ borel-measurable (F i) if  $j \leq i$  for j using that adaptedD by blast
  thus ( $\lambda\xi. \sum_{i < i} X (\text{Suc } i) \xi$ ) ∈ borel-measurable (F i) by auto
qed

lemma (in enat-filtered-measure) partial-sum-eSuc-adapted:
  assumes adapted-process M F 0 X
  shows adapted-process M F 0 ( $\lambda n \xi. \sum_{i < n} X (\text{eSuc } i) \xi$ )
proof (unfold-locales)
  interpret adapted-process M F 0 X using assms by blast
  fix i
  have X (eSuc j) ∈ borel-measurable (F i) if  $j < i$  for j using that adaptedD by
(simp add: ileI1)
  thus ( $\lambda\xi. \sum_{i < i} X (\text{eSuc } i) \xi$ ) ∈ borel-measurable (F i) by auto
qed

lemma (in filtered-measure) adapted-process-sum:
  assumes  $\bigwedge i. i \in I \implies$  adapted-process M F t₀ (X i)
  shows adapted-process M F t₀ ( $\lambda k \xi. \sum_{i \in I} X i k \xi$ )
proof -
  {
    fix i k assume i ∈ I and asm:  $t_0 \leq k$ 
    then interpret adapted-process M F t₀ X i using assms by simp
    have X i k ∈ borel-measurable M X i k ∈ borel-measurable (F k) using measurable-from-subalg subalgebras adapted asm by (blast, simp)
  }

```

```

thus ?thesis by (unfold-locales) simp
qed

```

An adapted process is necessarily a stochastic process.

```

sublocale adapted-process ⊆ stochastic-process using measurable-from-subalg sub-
algebras adapted by (unfold-locales) blast

```

A stochastic process is always adapted to the natural filtration it generates.

```

lemma (in stochastic-process) adapted-process-natural-filtration: adapted-process
M (natural-filtration M t₀ X) t₀ X
  using filtered-measure-natural-filtration
  by (intro-locales) (auto simp add: natural-filtration-def intro!: adapted-process-axioms.intro
measurable-family-vimage-algebra)

```

### 4.3 Progressively Measurable Process

```

locale progressive-process = filtered-measure M F t₀ for M F t₀ and X :: - ⇒ -
  ⇒ - :: {second-countable-topology, banach} +
  assumes progressive[measurable]: ∀ t. t₀ ≤ t ⇒ (λ(i, x). X i x) ∈ borel-measurable
  (restrict-space borel {t₀..t} ⊗ M F t)
begin

lemma progressiveD:
  assumes S ∈ borel
  shows (λ(j, ξ). X j ξ) −` S ∩ ({t₀..i} × space M) ∈ (restrict-space borel {t₀..i})
  ⊗ M F i)
  using measurable-sets[OF progressive, OF - assms, of i]
  by (cases t₀ ≤ i) (auto simp add: space-restrict-space sets-pair-measure space-pair-measure)

end

lemma (in filtered-measure) progressive-process-const-fun:
  assumes f ∈ borel-measurable (F t₀)
  shows progressive-process M F t₀ (λ-. f)
proof (unfold-locales)
  fix i assume asm: t₀ ≤ i
  have f ∈ borel-measurable (F i) using borel-measurable-mono[OF order.refl asm]
  assms by blast
  thus case-prod (λ-. f) ∈ borel-measurable (restrict-space borel {t₀..i} ⊗ M F i)
  using measurable-compose[OF measurable-snd] by simp
qed

lemma (in filtered-measure) progressive-process-const:
  assumes c ∈ borel-measurable borel
  shows progressive-process M F t₀ (λ i -. c i)
  using assms by (unfold-locales) (auto simp add: measurable-split-conv intro!
measurable-compose[OF measurable-fst] measurable-restrict-space1)

context progressive-process

```

```

begin

lemma compose-progressive:
  assumes case-prod f ∈ borel-measurable borel
  shows progressive-process M F t₀ (λi ξ. (f i) (X i ξ))
proof
  fix i assume asm: t₀ ≤ i
  have (λ(j, ξ). (j, X j ξ)) ∈ (restrict-space borel {t₀..i} ⊗ M F i) → M borel ⊗ M
  borel
  using progressive[OF asm] measurable-fst''[OF measurable-restrict-space1, OF
  measurable-id]
  by (auto simp add: measurable-pair-iff measurable-split-conv)
  moreover have (λ(j, ξ). f j (X j ξ)) = case-prod f o ((λ(j, y). (j, y)) o (λ(j, ξ).
  (j, X j ξ))) by fastforce
  ultimately show (λ(j, ξ). (f j) (X j ξ)) ∈ borel-measurable (restrict-space borel
  {t₀..i} ⊗ M F i) using assms by (simp add: borel-prod)
qed

lemma norm-progressive: progressive-process M F t₀ (λi ξ. norm (X i ξ)) using
measurable-compose[OF progressive borel-measurable-norm] by (unfold-locales)
simp

lemma scaleR-right-progressive:
  assumes progressive-process M F t₀ R
  shows progressive-process M F t₀ (λi ξ. (R i ξ) *ᵣ (X i ξ))
  using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)

lemma scaleR-right-const-fun-progressive:
  assumes f ∈ borel-measurable (F t₀)
  shows progressive-process M F t₀ (λi ξ. f ξ *ᵣ (X i ξ))
  using assms by (fast intro: scaleR-right-progressive progressive-process-const-fun)

lemma scaleR-right-const-progressive:
  assumes c ∈ borel-measurable borel
  shows progressive-process M F t₀ (λi ξ. c i *ᵣ (X i ξ))
  using assms by (fastforce intro: scaleR-right-progressive progressive-process-const)

lemma add-progressive:
  assumes progressive-process M F t₀ Y
  shows progressive-process M F t₀ (λi ξ. X i ξ + Y i ξ)
  using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)

lemma diff-progressive:
  assumes progressive-process M F t₀ Y
  shows progressive-process M F t₀ (λi ξ. X i ξ - Y i ξ)
  using progressive-process.progressive[OF assms] by (unfold-locales) (simp add:
progressive assms)

```

```

lemma uminus-progressive: progressive-process M F t0 (−X) using scaleR-right-const-progressive[of
λ-. −1] by (simp add: fun-Compl-def)

```

```

end

```

A progressively measurable process is also adapted.

```

sublocale progressive-process ⊆ adapted-process using measurable-compose-rev[OF
progressive measurable-Pair1]
  unfolding prod.case space-restrict-space
  by unfold-locales simp

```

In the discrete setting, adaptedness is equivalent to progressive measurability.

```

theorem (in nat-filtered-measure) progressive-iff-adapted: progressive-process M F
0 X ←→ adapted-process M F 0 X

```

```

proof (intro iffI)

```

```

  assume asm: progressive-process M F 0 X

```

```

  interpret progressive-process M F 0 X by (rule asm)

```

```

  show adapted-process M F 0 X ..

```

```

next

```

```

  assume asm: adapted-process M F 0 X

```

```

  interpret adapted-process M F 0 X by (rule asm)

```

```

  show progressive-process M F 0 X

```

```

  proof (unfold-locales, intro borel-measurableI)

```

```

    fix S :: 'b set and i :: nat assume open-S: open S
  {

```

```

    fix j assume asm: j ≤ i

```

```

    hence X j −‘ S ∩ space M ∈ F i using adaptedD[of j, THEN measurable-sets]
space-F open-S by fastforce

```

```

    moreover have case-prod X −‘ S ∩ {j} × space M = {j} × (X j −‘ S ∩
space M) for j by fast

```

```

    moreover have {j :: nat} ∈ restrict-space borel {0..i} using asm by (simp
add: sets-restrict-space-iff)

```

```

    ultimately have case-prod X −‘ S ∩ {j} × space M ∈ restrict-space borel
{0..i} ⊗M F i by simp
  }

```

```

    hence (λj. (λ(x, y). X x y) −‘ S ∩ {j} × space M) ‘ {..i} ⊆ restrict-space borel
{0..i} ⊗M F i by blast

```

```

    moreover have case-prod X −‘ S ∩ space (restrict-space borel {0..i} ⊗M F
i) = (⋃j≤i. case-prod X −‘ S ∩ {j} × space M) unfolding space-pair-measure
space-restrict-space space-F by force

```

```

    ultimately show case-prod X −‘ S ∩ space (restrict-space borel {0..i} ⊗M
F i) ∈ restrict-space borel {0..i} ⊗M F i by (metis sets.countable-UN)

```

```

qed

```

```

qed

```

```

theorem (in enat-filtered-measure) progressive-iff-adapted: progressive-process M
F 0 X ←→ adapted-process M F 0 X

```

```

proof (intro iffI)
  assume asm: progressive-process M F 0 X
  interpret progressive-process M F 0 X by (rule asm)
  show adapted-process M F 0 X ..
next
  assume asm: adapted-process M F 0 X
  interpret adapted-process M F 0 X by (rule asm)
  show progressive-process M F 0 X
  proof (unfold-locales, intro borel-measurableI)
    fix S :: 'b set and i :: enat assume open-S: open S
    {
      fix j assume asm: j ≤ i
      hence X j −‘ S ∩ space M ∈ F i using adaptedD[of j, THEN measurable-sets]
space-F open-S by fastforce
      moreover have case-prod X −‘ S ∩ {j} × space M = {j} × (X j −‘ S ∩
space M) for j by fast
      moreover have {j :: enat} ∈ restrict-space borel {0..i} using asm by (simp
add: sets-restrict-space-iff)
        ultimately have case-prod X −‘ S ∩ {j} × space M ∈ restrict-space borel
{0..i} ⊗M F i by simp
      }
      hence (λj. (λ(x, y). X x y) −‘ S ∩ {j} × space M) ‘{..i} ⊆ restrict-space borel
{0..i} ⊗M F i by blast
      moreover have case-prod X −‘ S ∩ space (restrict-space borel {0..i}) ⊗ M F i =
(∪j ≤ i. case-prod X −‘ S ∩ {j} × space M) unfolding space-pair-measure
space-restrict-space space-F by force
      ultimately show case-prod X −‘ S ∩ space (restrict-space borel {0..i}) ⊗ M F i ∈ restrict-space borel {0..i} ⊗M F i by (metis sets.countable-UN)
      qed
    qed

```

#### 4.4 Predictable Process

We introduce the constant  $\Sigma_P$  to denote the predictable  $\sigma$ -algebra.

```

context linearly-filtered-measure
begin

```

```

definition ΣP :: ('b × 'a) measure where predictable-sigma: ΣP ≡ sigma ({t0..} × space M) ({{s<..t}} × A | A s.t. A ∈ F s ∧ t0 ≤ s ∧ s < t} ∪ {{t0} × A | A ∈ F t0

```

```

lemma space-predictable-sigma[simp]: space ΣP = ({t0..} × space M) unfolding
predictable-sigma space-measure-of-conv by blast

```

```

lemma sets-predictable-sigma: sets ΣP = sigma-sets ({t0..} × space M) ({{s<..t}} × A | A ∈ F s ∧ t0 ≤ s ∧ s < t} ∪ {{t0} × A | A. A ∈ F t0})
unfolding predictable-sigma using space-F sets.sets-into-space by (subst sets-measure-of)
fastforce+

```

```

lemma measurable-predictable-sigma-snd:
  assumes countable  $\mathcal{I}$   $\mathcal{I} \subseteq \{\{s <.. t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0 <..\} \subseteq (\bigcup \mathcal{I})$ 
  shows  $\text{snd} \in \Sigma_P \rightarrow_M F t_0$ 
  proof (intro measurableI)
    fix  $S :: 'a set$  assume  $\text{asm}: S \in F t_0$ 
    have countable: countable  $((\lambda I. I \times S) ` \mathcal{I})$  using assms(1) by blast
    have  $(\lambda I. I \times S) ` \mathcal{I} \subseteq \{\{s <.. t\} \times A \mid A \text{ s t. } A \in F s \wedge t_0 \leq s \wedge s < t\}$  using
      sets-F-mono[ $\text{OF order-refl}$ , THEN subsetD,  $\text{OF - asm}$ ] assms(2) by blast
    hence  $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S \in \Sigma_P$  unfolding sets-predictable-sigma using
       $\text{asm by (intro sigma-sets-Un [OF sigma-sets-UNION [OF countable] sigma-sets.Basic] sigma-sets.Basic) blast+}$ 
    moreover have  $\text{snd} - ` S \cap \text{space } \Sigma_P = \{t_0..\} \times S$  using sets.sets-into-space[ $\text{OF}$ 
       $\text{asm}$ ] by fastforce
    moreover have  $\{t_0\} \cup \{t_0 <..\} = \{t_0..\}$  by auto
    moreover have  $(\bigcup I \in \mathcal{I}. I \times S) \cup \{t_0\} \times S = \{t_0..\} \times S$  using assms(2,3)
    calculation(3) by fastforce
    ultimately show  $\text{snd} - ` S \cap \text{space } \Sigma_P \in \Sigma_P$  by argo
  qed (auto)

lemma measurable-predictable-sigma-fst:
  assumes countable  $\mathcal{I}$   $\mathcal{I} \subseteq \{\{s <.. t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\} \{t_0 <..\} \subseteq (\bigcup \mathcal{I})$ 
  shows  $\text{fst} \in \Sigma_P \rightarrow_M \text{borel}$ 
  proof -
    have  $A \times \text{space } M \in \text{sets } \Sigma_P$  if  $A \in \text{sigma-sets } \{t_0..\} \{\{s <.. t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$  for  $A$  unfolding sets-predictable-sigma using that
    proof (induction rule: sigma-sets.induct)
      case (Basic a)
      thus ?case using space-F sets.top by blast
    next
      case (Compl a)
      have  $(\{t_0..\} - a) \times \text{space } M = \{t_0..\} \times \text{space } M - a \times \text{space } M$  by blast
      then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
    next
      case (Union a)
      have  $\bigcup (\text{range } a) \times \text{space } M = \bigcup (\text{range } (\lambda i. a \ i \times \text{space } M))$  by blast
      then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
    qed (auto)
    moreover have restrict-space borel  $\{t_0..\} = \text{sigma } \{t_0..\} \{\{s <.. t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ 
    proof -
      have sigma-sets  $\{t_0..\} ((\cap) \{t_0..\} ` \text{sigma-sets UNIV} (\text{range greaterThan})) =$ 
        sigma-sets  $\{t_0..\} \{\{s <.. t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ 
      proof (intro sigma-sets-eqI ; clarify)
        fix  $A :: 'b set$  assume  $\text{asm}: A \in \text{sigma-sets UNIV} (\text{range greaterThan})$ 
        thus  $\{t_0..\} \cap A \in \text{sigma-sets } \{t_0..\} \{\{s <.. t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ 
        proof (induction rule: sigma-sets.induct)
          case (Basic a)
          then obtain  $s$  where  $s: a = \{s <..\}$  by blast
          show ?case

```

```

proof (cases  $t_0 \leq s$ )
  case True
    hence  $\{t_0..\} \cap a = (\bigcup i \in \mathcal{I}. \{s<.. \} \cap i)$  using  $s$  assms(3) by force
    have  $((\cap) \{s<.. \} \cdot \mathcal{I}) \subseteq \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$ 
    proof (clarify)
      fix  $A$  assume  $A \in \mathcal{I}$ 
      then obtain  $s' t'$  where  $A: A = \{s'<..t'\} t_0 \leq s' s' < t'$  using assms(2)
    by blast
      hence  $\{s<.. \} \cap A = \{\max s s' <.. t'\}$  by fastforce
      moreover have  $t_0 \leq \max s s'$  using  $A$  True by linarith
      moreover have  $\max s s' < t'$  if  $s < t'$  using  $A$  that by linarith
      moreover have  $\{s<.. \} \cap A = \{\}$  if  $\neg s < t'$  using  $A$  that by force
      ultimately show  $\{s<.. \} \cap A \in \text{sigma-sets } \{t_0..\} \{\{s<..t\} \mid s \text{ t. } t_0 \leq s \wedge s < t\}$  by (cases  $s < t'$ ) (blast, simp add: sigma-sets.Empty)
    qed
    thus ?thesis unfolding * using assms(1) by (intro sigma-sets-UNION)
  auto
  next
    case False
    hence  $\{t_0..\} \cap a = \{t_0..\}$  using  $s$  by force
    thus ?thesis using sigma-sets-top by auto
  qed
  next
    case (Compl  $a$ )
    have  $\{t_0..\} \cap (UNIV - a) = \{t_0..\} - (\{t_0..\} \cap a)$  by blast
    then show ?case using Compl(2)[THEN sigma-sets.Compl] by presburger
  next
    case (Union  $a$ )
    have  $\{t_0..\} \cap \bigcup (\text{range } a) = \bigcup (\text{range } (\lambda i. \{t_0..\} \cap a_i))$  by blast
    then show ?case using Union(2)[THEN sigma-sets.Union] by presburger
  qed (simp add: sigma-sets.Empty)
  next
    fix  $s t$  assume  $asm: t_0 \leq s s < t$ 
    hence  $\{s<..t\} = \{s<.. \} \cap (\{t_0..\} - \{t<.. \})$  by force
    have  $\{s<.. \} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \cdot \text{sigma-sets } UNIV (\text{range greaterThan}))$ 
    using  $asm$  by (intro sigma-sets.Basic) auto
    moreover have  $\{t_0..\} - \{t<.. \} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \cdot \text{sigma-sets } UNIV (\text{range greaterThan}))$  using  $asm$  by (intro sigma-sets.Compl sigma-sets.Basic)
    auto
    ultimately show  $\{s<..t\} \in \text{sigma-sets } \{t_0..\} ((\cap) \{t_0..\} \cdot \text{sigma-sets } UNIV (\text{range greaterThan}))$  unfolding * Int-range-binary[of  $\{s<.. \}$ ] by (intro sigma-sets-Inter[OF - binary-in-sigma-sets]) auto
  qed
  thus ?thesis unfolding borel-Ioi restrict-space-def emeasure-sigma by (force intro: sigma-eqI)
  qed
  ultimately have restrict-space borel  $\{t_0..\} \otimes_M \text{sigma } (\text{space } M) \{\} \subseteq \text{sets } \Sigma_P$ 
  unfolding sets-pair-measure space-restrict-space space-measure-of-conv

```

```

using space-predictable-sigma sets.sigma-algebra-axioms[of  $\Sigma_P$ ]
by (intro sigma-algebra.sigma-sets-subset) (auto simp add: sigma-sets-empty-eq
sets-measure-of-conv)
moreover have space (restrict-space borel {t0..}  $\otimes_M$  sigma (space M) {}) =
space  $\Sigma_P$  by (simp add: space-pair-measure)
moreover have fst  $\in$  restrict-space borel {t0..}  $\otimes_M$  sigma (space M) {}  $\rightarrow_M$ 
borel by (fastforce intro: measurable-fst'[OF measurable-restrict-space1, of  $\lambda x. x$ ])

ultimately show ?thesis by (meson borel-measurable-subalgebra)
qed

end

locale predictable-process = linearly-filtered-measure M F t0 for M F t0 and X :: 
-  $\Rightarrow$  -  $\Rightarrow$  - :: {second-countable-topology, banach} +
assumes predictable:  $(\lambda(t, x). X t x) \in$  borel-measurable  $\Sigma_P$ 
begin

lemmas predictableD = measurable-sets[OF predictable, unfolded space-predictable-sigma]

end

lemma (in nat-filtered-measure) measurable-predictable-sigma-snd':
shows snd  $\in$   $\Sigma_P \rightarrow_M F 0$ 
by (intro measurable-predictable-sigma-snd[of range  $(\lambda x. \{Suc x\})$ ]) (force | simp
add: greaterThan-0)+

lemma (in nat-filtered-measure) measurable-predictable-sigma-fst':
shows fst  $\in$   $\Sigma_P \rightarrow_M$  borel
by (intro measurable-predictable-sigma-fst[of range  $(\lambda x. \{Suc x\})$ ]) (force | simp
add: greaterThan-0)+

lemma (in enat-filtered-measure) measurable-predictable-sigma-snd':
shows snd  $\in$   $\Sigma_P \rightarrow_M F 0$ 
by (intro measurable-predictable-sigma-snd[of  $\{\{0 <..\infty\}\}$ ]) force+

lemma (in enat-filtered-measure) measurable-predictable-sigma-fst':
shows fst  $\in$   $\Sigma_P \rightarrow_M$  borel
by (intro measurable-predictable-sigma-fst[of  $\{\{0 <..\infty\}\}$ ]) force+

lemma (in real-filtered-measure) measurable-predictable-sigma-snd':
shows snd  $\in$   $\Sigma_P \rightarrow_M F 0$ 
using real-arch-simple by (intro measurable-predictable-sigma-snd[of range  $(\lambda x::nat. \{0 <..\text{real } (Suc x)\})$ ]) (fastforce intro: add-increasing)+

lemma (in real-filtered-measure) measurable-predictable-sigma-fst':
shows fst  $\in$   $\Sigma_P \rightarrow_M$  borel
using real-arch-simple by (intro measurable-predictable-sigma-fst[of range  $(\lambda x::nat. \{0 <..\text{real } (Suc x)\})$ ]) (fastforce intro: add-increasing)+
```

```

lemma (in ennreal-filtered-measure) measurable-predictable-sigma-snd':
  shows snd ∈ ΣP →M F 0
  by (intro measurable-predictable-sigma-snd[of {{0<..∞}}]) force+
lemma (in ennreal-filtered-measure) measurable-predictable-sigma-fst':
  shows fst ∈ ΣP →M borel
  by (intro measurable-predictable-sigma-fst[of {{0<..∞}}]) force+

We show sufficient conditions for functions constant in one argument to
constitute a predictable process. In contrast to the cases before, this is not
a triviality.

lemma (in linearly-filtered-measure) predictable-process-const-fun:
  assumes snd ∈ ΣP →M F t0 f ∈ borel-measurable (F t0)
  shows predictable-process M F t0 (λ-. f)
  using measurable-compose-rev[OF assms(2)] assms(1) by (unfold-locales) (auto
simp add: measurable-split-conv)

lemma (in nat-filtered-measure) predictable-process-const-fun'[intro]:
  assumes f ∈ borel-measurable (F 0)
  shows predictable-process M F 0 (λ-. f)
  using assms by (intro predictable-process-const-fun[OF measurable-predictable-sigma-snd'])

lemma (in enat-filtered-measure) predictable-process-const-fun'[intro]:
  assumes f ∈ borel-measurable (F 0)
  shows predictable-process M F 0 (λ-. f)
  using assms by (intro predictable-process-const-fun[OF measurable-predictable-sigma-snd'])

lemma (in real-filtered-measure) predictable-process-const-fun'[intro]:
  assumes f ∈ borel-measurable (F 0)
  shows predictable-process M F 0 (λ-. f)
  using assms by (intro predictable-process-const-fun[OF measurable-predictable-sigma-snd'])

lemma (in ennreal-filtered-measure) predictable-process-const-fun'[intro]:
  assumes f ∈ borel-measurable (F 0)
  shows predictable-process M F 0 (λ-. f)
  using assms by (intro predictable-process-const-fun[OF measurable-predictable-sigma-snd'])

lemma (in linearly-filtered-measure) predictable-process-const:
  assumes fst ∈ borel-measurable ΣP c ∈ borel-measurable borel
  shows predictable-process M F t0 (λi -. c i)
  using assms by (unfold-locales) (simp add: measurable-split-conv)

lemma (in linearly-filtered-measure) predictable-process-const-const[intro]:
  shows predictable-process M F t0 (λ- -. c)
  by (unfold-locales) simp

lemma (in nat-filtered-measure) predictable-process-const'[intro]:
  assumes c ∈ borel-measurable borel

```

```

shows predictable-process M F 0 ( $\lambda i \_. c i$ )
using assms by (intro predictable-process-const[OF measurable-predictable-sigma-fst])

lemma (in enat-filtered-measure) predictable-process-const'[intro]:
assumes  $c \in \text{borel-measurable borel}$ 
shows predictable-process M F 0 ( $\lambda i \_. c i$ )
using assms by (intro predictable-process-const[OF measurable-predictable-sigma-fst])

lemma (in real-filtered-measure) predictable-process-const'[intro]:
assumes  $c \in \text{borel-measurable borel}$ 
shows predictable-process M F 0 ( $\lambda i \_. c i$ )
using assms by (intro predictable-process-const[OF measurable-predictable-sigma-fst])

lemma (in ennreal-filtered-measure) predictable-process-const'[intro]:
assumes  $c \in \text{borel-measurable borel}$ 
shows predictable-process M F 0 ( $\lambda i \_. c i$ )
using assms by (intro predictable-process-const[OF measurable-predictable-sigma-fst])

context predictable-process
begin

lemma compose-predictable:
assumes  $\text{fst} \in \text{borel-measurable } \Sigma_P \text{ case-prod } f \in \text{borel-measurable borel}$ 
shows predictable-process M F  $t_0$  ( $\lambda i \xi. (f i) (X i \xi)$ )
proof
have  $(\lambda(i, \xi). (i, X i \xi)) \in \Sigma_P \rightarrow_M \text{borel} \otimes_M \text{borel}$  using predictable assms(1)
by (auto simp add: measurable-pair-iff measurable-split-conv)
moreover have  $(\lambda(i, \xi). f i (X i \xi)) = \text{case-prod } f o (\lambda(i, \xi). (i, X i \xi))$  by
fastforce
ultimately show  $(\lambda(i, \xi). f i (X i \xi)) \in \text{borel-measurable } \Sigma_P$  unfolding borel-prod
using assms by simp
qed

lemma norm-predictable: predictable-process M F  $t_0$  ( $\lambda i \xi. \text{norm } (X i \xi)$ ) using
measurable-compose[OF predictable borel-measurable-norm]
by (unfold-locales) (simp add: prod.case-distrib)

lemma scaleR-right-predictable:
assumes predictable-process M F  $t_0 R$ 
shows predictable-process M F  $t_0$  ( $\lambda i \xi. (R i \xi) *_R (X i \xi)$ )
using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)

lemma scaleR-right-const-fun-predictable:
assumes  $\text{snd} \in \Sigma_P \rightarrow_M F t_0 f \in \text{borel-measurable } (F t_0)$ 
shows predictable-process M F  $t_0$  ( $\lambda i \xi. f \xi *_R (X i \xi)$ )
using assms by (fast intro: scaleR-right-predictable predictable-process-const-fun)

lemma scaleR-right-const-predictable:

```

```

assumes fst ∈ borel-measurable ΣP c ∈ borel-measurable borel
shows predictable-process M F t0 (λi ξ. c i *R (X i ξ))
using assms by (fastforce intro: scaleR-right-predictable predictable-process-const)

lemma scaleR-right-const'-predictable: predictable-process M F t0 (λi ξ. c *R (X i ξ))
by (fastforce intro: scaleR-right-predictable)

lemma add-predictable:
assumes predictable-process M F t0 Y
shows predictable-process M F t0 (λi ξ. X i ξ + Y i ξ)
using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)

lemma diff-predictable:
assumes predictable-process M F t0 Y
shows predictable-process M F t0 (λi ξ. X i ξ - Y i ξ)
using predictable predictable-process.predictable[OF assms] by (unfold-locales)
(auto simp add: measurable-split-conv)

lemma uminus-predictable: predictable-process M F t0 (-X) using scaleR-right-const'-predictable[of -1]
by (simp add: fun-Compl-def)

end

Every predictable process is also progressively measurable.

sublocale predictable-process ⊆ progressive-process
proof (unfold-locales)
fix i :: 'b assume asm: t0 ≤ i
{
  fix S :: ('b × 'a) set assume S ∈ {{s<..t} × A | A s t. A ∈ F s ∧ t0 ≤ s ∧ s < t} ∪ {{t0} × A | A. A ∈ F t0}
  hence (λx. x) -` S ∩ ({t0..i} × space M) ∈ restrict-space borel {t0..i} ⊗M F
i
  proof
    assume S ∈ {{s<..t} × A | A s t. A ∈ F s ∧ t0 ≤ s ∧ s < t}
    then obtain s t A where S-is: S = {s<..t} × A t0 ≤ s s < t A ∈ F s by blast
    hence (λx. x) -` S ∩ ({t0..i} × space M) = {s<..min i t} × A using sets.sets-into-space[OF S-is(4)] by auto
    then show ?thesis using S-is sets-F-mono[of s i] by (cases s ≤ i) (fastforce simp add: sets-restrict-space-iff) +
  next
    assume S ∈ {{t0} × A | A. A ∈ F t0}
    then obtain A where S-is: S = {t0} × A A ∈ F t0 by blast
    hence (λx. x) -` S ∩ ({t0..i} × space M) = {t0} × A using asm sets.sets-into-space[OF S-is(2)] by auto
    thus ?thesis using S-is(2) sets-F-mono[OF order-refl asm] asm by (fastforce simp add: sets-restrict-space-iff)

```

```

qed
hence  $(\lambda x. x) -` S \cap space (restrict-space borel \{t_0..i\} \otimes_M F i) \in restrict-space$ 
borel  $\{t_0..i\} \otimes_M F i$  by (simp add: space-pair-measure space-F[OF asm])
}
moreover have  $\{\{s <.. t\} \times A \mid A \in sets (F s) \wedge t_0 \leq s \wedge s < t\} \cup \{\{t_0\}$ 
 $\times A \mid A \in sets (F t_0)\} \subseteq Pow (\{t_0..\} \times space M)$  using sets.sets-into-space by
force
ultimately have  $(\lambda x. x) \in restrict-space borel \{t_0..i\} \otimes_M F i \rightarrow_M \Sigma_P$  us-
ing space-F[OF asm] by (intro measurable-sigma-sets[OF sets-predictable-sigma])
(fast, force simp add: space-pair-measure)
thus case-prod  $X \in borel-measurable (restrict-space borel \{t_0..i\} \otimes_M F i)$  using
predictable by simp
qed

```

The following lemma characterizes predictability in a discrete-time setting.

```

lemma (in nat-filtered-measure) sets-in-filtration:
assumes  $(\bigcup i. \{i\} \times A i) \in \Sigma_P$ 
shows  $A (Suc i) \in F i$   $A 0 \in F 0$ 
using assms unfolding sets-predictable-sigma
proof (induction  $(\bigcup i. \{i\} \times A i)$  arbitrary:  $A$ )
case Basic
{
assume  $\exists S. (\bigcup i. \{i\} \times A i) = \{0\} \times S$ 
then obtain  $S$  where  $S: (\bigcup i. \{i\} \times A i) = \{0\} \times S$  by blast
hence  $S \in F 0$  using Basic by (fastforce simp add: times-eq-iff)
moreover have  $A i = \{\}$  if  $i \neq 0$  for  $i$  using that  $S$  unfolding bot-nat-def[symmetric]
by blast
moreover have  $A 0 = S$  using  $S$  by blast
ultimately have  $A 0 \in F 0$   $A (Suc i) \in F i$  for  $i$  by auto
}
note * = this
{
assume  $\nexists S. (\bigcup i. \{i\} \times A i) = \{0\} \times S$ 
then obtain  $s t B$  where  $B: (\bigcup i. \{i\} \times A i) = \{s <.. t\} \times B$   $B \in sets (F s)$ 
 $s < t$  using Basic by auto
hence  $A i = B$  if  $i \in \{s <.. t\}$  for  $i$  using that by fast
moreover have  $A i = \{\}$  if  $i \notin \{s <.. t\}$  for  $i$  using  $B$  that by fastforce
ultimately have  $A 0 \in F 0$   $A (Suc i) \in F i$  for  $i$  using  $B$  sets-F-mono
by (simp, metis less-Suc-eq-le sets.empty-sets subset-eq bot-nat-0.extremum
greaterThanAtMost-iff)
}
note ** = this
show  $A (Suc i) \in sets (F i)$   $A 0 \in F 0$  using *(2)[of i] *(1) **(2)[of i] **(1)
by blast+
next
case Empty
{
case 1

```

```

    then show ?case using Empty by simp
next
  case 2
    then show ?case using Empty by simp
  }
next
  case (Compl a)
  have a-in:  $a \subseteq \{0..\} \times \text{space } M$  using Compl(1) sets.sets-into-space sets-predictable-sigma
  space-predictable-sigma by metis
  hence A-in:  $A i \subseteq \text{space } M$  for  $i$  using Compl(4) by blast
  have a:  $a = \{0..\} \times \text{space } M - (\bigcup i. \{i\} \times A i)$  using a-in Compl(4) by blast
  also have ... =  $(\bigcap j. -(\{j\} \times (\text{space } M - A j)))$  by blast
  also have ... =  $(\bigcup j. \{j\} \times (\text{space } M - A j))$  by blast
  finally have *:  $(\text{space } M - A (\text{Suc } i)) \in F i$   $(\text{space } M - A 0) \in F 0$  using
  Compl(2,3) by auto
  {
    case 1
      then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
  sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
    next
      case 2
        then show ?case using * A-in by (metis bot-nat-0.extremum double-diff
  sets.Diff sets.top sets-F-mono sets-le-imp-space-le space-F)
    }
  next
  case (Union a)
  have a-in:  $a i \subseteq \{0..\} \times \text{space } M$  for  $i$  using Union(1) sets.sets-into-space
  sets-predictable-sigma space-predictable-sigma by metis
  hence A-in:  $A i \subseteq \text{space } M$  for  $i$  using Union(4) by blast
  have snd x ∈ snd ‘( $a i \cap (\{\text{fst } x\} \times \text{space } M)$ ) if  $x \in a i$  for  $i$   $x$  using that
  a-in by fastforce
  hence a-i:  $a i = (\bigcup j. \{j\} \times (\text{snd } ' (a i \cap (\{j\} \times \text{space } M))))$  for  $i$  by force
  have A-i:  $A i = \text{snd } ' (\bigcup (\text{range } a) \cap (\{i\} \times \text{space } M))$  for  $i$  unfolding
  Union(4) using A-in by force
  have *:  $\text{snd } ' (a j \cap (\{\text{Suc } i\} \times \text{space } M)) \in F i$   $\text{snd } ' (a j \cap (\{0\} \times \text{space } M))$ 
  ∈ F 0 for  $j$  using Union(2,3)[OF a-i] by auto
  {
    case 1
      have  $(\bigcup j. \text{snd } ' (a j \cap (\{\text{Suc } i\} \times \text{space } M))) \in F i$  using * by fast
      moreover have  $(\bigcup j. \text{snd } ' (a j \cap (\{\text{Suc } i\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap (\{\text{Suc } i\} \times \text{space } M))$  by fast
      ultimately show ?case using A-i by metis
    next
      case 2
        have  $(\bigcup j. \text{snd } ' (a j \cap (\{0\} \times \text{space } M))) \in F 0$  using * by fast
        moreover have  $(\bigcup j. \text{snd } ' (a j \cap (\{0\} \times \text{space } M))) = \text{snd } ' (\bigcup (\text{range } a) \cap (\{0\} \times \text{space } M))$  by fast
        ultimately show ?case using A-i by metis
  }

```

**qed**

This leads to the following useful fact.

```

lemma (in nat-filtered-measure) predictable-implies-adapted-Suc:
  assumes predictable-process M F 0 X
  shows adapted-process M F 0 ( $\lambda i. X(Suc i)$ )
proof (unfold-locales, intro borel-measurableI)
  interpret predictable-process M F 0 X by (rule assms)
  fix S :: 'b set and i assume open-S: open S
  have {Suc i} = {i <.. Suc i} by fastforce
  hence {Suc i} × space M ∈  $\Sigma_P$  using space-F[symmetric, of i] unfolding
  sets-predictable-sigma by (intro sigma-sets.Basic) blast
  moreover have case-prod X -' S ∩ (UNIV × space M) ∈  $\Sigma_P$  unfolding
  atLeast-0[symmetric] using open-S by (intro predictableD, simp add: borel-open)
  ultimately have case-prod X -' S ∩ ({Suc i} × space M) ∈  $\Sigma_P$  unfolding
  sets-predictable-sigma using space-F sets.sets-into-space
  by (subst Times-Int-distrib1[of {Suc i} UNIV space M, simplified], subst
  inf.commute, subst Int-assoc[symmetric], subst Int-range-binary)
  (intro sigma-sets-Inter binary-in-sigma-sets, fast)+
  moreover have case-prod X -' S ∩ ({Suc i} × space M) = {Suc i} × (X (Suc
  i) -' S ∩ space M) by (auto simp add: le-Suc-eq)
  moreover have ... = ( $\bigcup j. \{j\} \times (\text{if } j = Suc i \text{ then } (X (Suc i) -' S \cap space M)$ 
  else {})) by (force split: if-splits)
  ultimately have ( $\bigcup j. \{j\} \times (\text{if } j = Suc i \text{ then } (X (Suc i) -' S \cap space M)$ 
  else {})) ∈  $\Sigma_P$  by argo
  thus X (Suc i) -' S ∩ space (F i) ∈ sets (F i) using sets-in-filtration[of λj.
  if j = Suc i then (X (Suc i) -' S ∩ space M) else {}] space-F[OF zero-le] by
  presburger
qed

```

The following lemma characterizes predictability in the discrete setting.

```

theorem (in nat-filtered-measure) predictable-process-iff: predictable-process M F
0 X  $\longleftrightarrow$  adapted-process M F 0 ( $\lambda i. X(Suc i)$ )  $\wedge$  X 0 ∈ borel-measurable (F 0)
proof (intro iffI)
  assume asm: adapted-process M F 0 ( $\lambda i. X(Suc i)$ )  $\wedge$  X 0 ∈ borel-measurable
  (F 0)
  interpret adapted-process M F 0  $\lambda i. X(Suc i)$  using asm by blast
  have ( $\lambda(x, y). X x y$ ) ∈ borel-measurable  $\Sigma_P$ 
proof (intro borel-measurableI)
  fix S :: 'b set assume open-S: open S
  have {i} × (X i -' S ∩ space M) ∈ sets  $\Sigma_P$  for i
  proof (cases i)
    case 0
    then show ?thesis unfolding sets-predictable-sigma
    using measurable-sets[OF - borel-open[OF open-S], of X 0 F 0] asm by auto
  next
    case (Suc i)
    have {Suc i} = {i <.. Suc i} by fastforce
    then show ?thesis unfolding sets-predictable-sigma

```

```

using measurable-sets[OF adapted borel-open[OF open-S], of i]
by (intro sigma-sets.Basic, auto simp add: Suc)
qed
moreover have  $(\lambda(x, y). X x y) -` S \cap space \Sigma_P = (\bigcup i. \{i\} \times (X i -` S \cap space M))$  by fastforce
ultimately show  $(\lambda(x, y). X x y) -` S \cap space \Sigma_P \in sets \Sigma_P$  by simp
qed
thus predictable-process M F 0 X by (unfold-locales)
next
assume asm: predictable-process M F 0 X
interpret predictable-process M F 0 X using asm by blast
show adapted-process M F 0  $(\lambda i. X (Suc i)) \wedge X 0 \in borel-measurable (F 0)$ 
using predictable-implies-adapted-Suc asm by auto
qed

corollary (in nat-filtered-measure) predictable-processI[intro!]:
assumes  $X 0 \in borel-measurable (F 0) \wedge \forall i. X (Suc i) \in borel-measurable (F i)$ 
shows predictable-process M F 0 X
unfolding predictable-process-iff
using assms
by (meson adapted-process.intro adapted-process-axioms-def filtered-measure-axioms)

end

```

```

theory Martingale
imports Stochastic-Process Conditional-Expectation-Banach
begin

```

## 5 Martingales

The following locales are necessary for defining martingales.

### 5.1 Martingale

A martingale is an adapted process where the expected value of the next observation, given all past observations, is equal to the current value.

```

locale martingale = sigma-finite-filtered-measure + adapted-process +
assumes integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M (X i)$ 
and martingale-property:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies AE \xi \text{ in } M. X i \xi = cond-exp M (F i) (X j) \xi$ 

locale martingale-order = martingale M F t0 X for M F t0 and X :: - ⇒ - ⇒ -
:: {order-topology, ordered-real-vector}
locale martingale-linorder = martingale M F t0 X for M F t0 and X :: - ⇒ - ⇒ -
:: {linorder-topology, ordered-real-vector}
sublocale martingale-linorder ⊆ martingale-order ..

```

```

lemma (in sigma-finite-filtered-measure) martingale-const-fun[intro]:
  assumes integrable M f f ∈ borel-measurable (F t₀)
  shows martingale M F t₀ (λ-. f)
  using assms sigma-finite-subalgebra.cond-exp-F-meas[OF - assms(1), THEN AE-symmetric]
borel-measurable-mono
  by (unfold-locales) blast+

lemma (in sigma-finite-filtered-measure) martingale-cond-exp[intro]:
  assumes integrable M f
  shows martingale M F t₀ (λi. cond-exp M (F i) f)
  using sigma-finite-subalgebra.borel-measurable-cond-exp' borel-measurable-cond-exp

  by (unfold-locales) (auto intro: sigma-finite-subalgebra.cond-exp-nested-subalg[OF
- assms] simp add: subalgebra-F subalgebras)

corollary (in sigma-finite-filtered-measure) martingale-zero[intro]: martingale M
F t₀ (λ- -. 0) by fastforce

corollary (in finite-filtered-measure) martingale-const[intro]: martingale M F t₀
(λ- -. c) by fastforce

```

## 5.2 Submartingale

A submartingale is an adapted process where the expected value of the next observation, given all past observations, is greater than or equal to the current value.

```

locale submartingale = sigma-finite-filtered-measure M F t₀ + adapted-process M
F t₀ X for M F t₀ and X :: - ⇒ - ⇒ - :: {order-topology, ordered-real-vector} +
assumes integrable: ∀i. t₀ ≤ i ⇒ integrable M (X i)
  and submartingale-property: ∀i j. t₀ ≤ i ⇒ i ≤ j ⇒ AE ξ in M. X i ξ ≤
cond-exp M (F i) (X j) ξ

```

```

locale submartingale-linorder = submartingale M F t₀ X for M F t₀ and X :: -
⇒ - ⇒ - :: {linorder-topology}

```

```

lemma (in sigma-finite-filtered-measure) submartingale-const-fun[intro]:
  assumes integrable M f f ∈ borel-measurable (F t₀)
  shows submartingale M F t₀ (λ-. f)
proof –
  interpret martingale M F t₀ λ-. f using assms by (rule martingale-const-fun)
  show submartingale M F t₀ (λ-. f) using martingale-property by (unfold-locales)
  (force simp add: integrable)+
qed

```

```

lemma (in sigma-finite-filtered-measure) submartingale-cond-exp[intro]:
  assumes integrable M f
  shows submartingale M F t₀ (λi. cond-exp M (F i) f)
proof –

```

```

interpret martingale M F t₀ λi. cond-exp M (F i) f using assms by (rule
martingale-cond-exp)
show submartingale M F t₀ (λi. cond-exp M (F i) f) using martingale-property
by (unfold-locales) (force simp add: integrable) +
qed

corollary (in finite-filtered-measure) submartingale-const[intro]: submartingale M
F t₀ (λ- -. c) by fastforce

sublocale martingale-order ⊆ submartingale using martingale-property by (unfold-locales)
(force simp add: integrable) +
sublocale martingale-linorder ⊆ submartingale-linorder ..

```

### 5.3 Supermartingale

A supermartingale is an adapted process where the expected value of the next observation, given all past observations, is less than or equal to the current value.

```

locale supermartingale = sigma-finite-filtered-measure M F t₀ + adapted-process
M F t₀ X for M F t₀ and X :: - ⇒ - ⇒ - :: {order-topology, ordered-real-vector}
+
assumes integrable: ∀i. t₀ ≤ i ⇒ integrable M (X i)
and supermartingale-property: ∀i j. t₀ ≤ i ⇒ i ≤ j ⇒ AE ξ in M. X i ξ
≥ cond-exp M (F i) (X j) ξ

locale supermartingale-linorder = supermartingale M F t₀ X for M F t₀ and X
:: - ⇒ - ⇒ - :: {linorder-topology}

```

```

lemma (in sigma-finite-filtered-measure) supermartingale-const-fun[intro]:
assumes integrable M f f ∈ borel-measurable (F t₀)
shows supermartingale M F t₀ (λ-. f)
proof -
interpret martingale M F t₀ λ-. f using assms by (rule martingale-const-fun)
show supermartingale M F t₀ (λ-. f) using martingale-property by (unfold-locales)
(force simp add: integrable) +
qed

```

```

lemma (in sigma-finite-filtered-measure) supermartingale-cond-exp[intro]:
assumes integrable M f
shows supermartingale M F t₀ (λi. cond-exp M (F i) f)
proof -
interpret martingale M F t₀ λi. cond-exp M (F i) f using assms by (rule
martingale-cond-exp)
show supermartingale M F t₀ (λi. cond-exp M (F i) f) using martingale-property
by (unfold-locales) (force simp add: integrable) +
qed

```

```

corollary (in finite-filtered-measure) supermartingale-const[intro]: supermartingale
M F t₀ (λ- -. c) by fastforce

```

```

sublocale martingale-order ⊆ supermartingale using martingale-property by (unfold-locales)
  (force simp add: integrable)+
sublocale martingale-linorder ⊆ supermartingale-linorder ..

```

A stochastic process is a martingale, if and only if it is both a submartingale and a supermartingale.

```

lemma martingale-iff:
  shows martingale M F t₀ X  $\longleftrightarrow$  submartingale M F t₀ X  $\wedge$  supermartingale M
  F t₀ X
proof (rule iffI)
  assume asm: martingale M F t₀ X
  interpret martingale-order M F t₀ X by (intro martingale-order.intro asm)
  show submartingale M F t₀ X  $\wedge$  supermartingale M F t₀ X using submartin-
  gale-axioms supermartingale-axioms by blast
next
  assume asm: submartingale M F t₀ X  $\wedge$  supermartingale M F t₀ X
  interpret submartingale M F t₀ X by (simp add: asm)
  interpret supermartingale M F t₀ X by (simp add: asm)
  show martingale M F t₀ X using submartingale-property supermartingale-property
  by (unfold-locales) (intro integrable, blast, force)
qed

```

## 5.4 Martingale Lemmas

In the following segment, we cover basic properties of martingales.

```

context martingale
begin

```

```

lemma cond-exp-diff-eq-zero:
  assumes t₀ ≤ i i ≤ j
  shows AE ξ in M. cond-exp M (F i) (λξ. X j ξ − X i ξ) ξ = 0
  using martingale-property[OF assms] assms
    sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i]
    sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i j i] by
    fastforce

lemma set-integral-eq:
  assumes A ∈ F i t₀ ≤ i i ≤ j
  shows set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j)
proof –
  interpret sigma-finite-subalgebra M F i using assms(2) by blast
  have (ʃ x ∈ A. X i x ∂M) = (ʃ x ∈ A. cond-exp M (F i) (X j) x ∂M) using
  martingale-property[OF assms(2,3)] borel-measurable-cond-exp' assms subalgebras
  subalgebra-def by (intro set-lebesgue-integral-cong-AE[OF - random-variable]) fast-
  force+
  also have ... = (ʃ x ∈ A. X j x ∂M) using assms by (auto simp: integrable
  intro: cond-exp-set-integral[symmetric])

```

**finally show** ?thesis .  
**qed**

**lemma** scaleR-const[intro]:  
**shows** martingale M F t0 ( $\lambda i x. c *_R X i x$ )  
**proof –**  
{  
fix i j :: 'b **assume** asm:  $t_0 \leq i \leq j$   
**interpret** sigma-finite-subalgebra M F i **using** asm **by** blast  
**have** AE x in M.  $c *_R X i x = cond-exp M (F i) (\lambda x. c *_R X j x) x$  **using** asm cond-exp-scaleR-right[OF integrable, of j, THEN AE-symmetric] martingale-property[OF asm] **by** force  
}  
**thus** ?thesis **by** (unfold-locales) (auto simp add: integrable martingale.integrable)  
**qed**

**lemma** uminus[intro]:  
**shows** martingale M F t0 ( $- X$ )  
**using** scaleR-const[of -1] **by** (force intro: back-subst[of martingale M F t0])

**lemma** add[intro]:  
**assumes** martingale M F t0 Y  
**shows** martingale M F t0 ( $\lambda i \xi. X i \xi + Y i \xi$ )  
**proof –**  
**interpret** Y: martingale M F t0 Y **by** (rule assms)  
{  
fix i j :: 'b **assume** asm:  $t_0 \leq i \leq j$   
**hence** AE  $\xi$  in M.  $X i \xi + Y i \xi = cond-exp M (F i) (\lambda x. X j x + Y j x) \xi$   
**using** sigma-finite-subalgebra.cond-exp-add[OF - integrable martingale.integrable[OF assms], of F i j j, THEN AE-symmetric]  
**martingale-property**[OF asm] martingale.martingale-property[OF assms  
asm] **by** force  
}  
**thus** ?thesis **using** assms  
**by** (unfold-locales) (auto simp add: integrable martingale.integrable)  
**qed**

**lemma** diff[intro]:  
**assumes** martingale M F t0 Y  
**shows** martingale M F t0 ( $\lambda i x. X i x - Y i x$ )  
**proof –**  
**interpret** Y: martingale M F t0 Y **by** (rule assms)  
{  
fix i j :: 'b **assume** asm:  $t_0 \leq i \leq j$   
**hence** AE  $\xi$  in M.  $X i \xi - Y i \xi = cond-exp M (F i) (\lambda x. X j x - Y j x) \xi$   
**using** sigma-finite-subalgebra.cond-exp-diff[OF - integrable martingale.integrable[OF assms], of F i j j, THEN AE-symmetric]  
**martingale-property**[OF asm] martingale.martingale-property[OF assms  
asm] **by** fastforce

```

}

thus ?thesis using assms by (unfold-locales) (auto simp add: integrable martingale.integrable)
qed

end

Using properties of the conditional expectation, we present the following alternative characterizations of martingales.

lemma (in sigma-finite-filtered-measure) martingale-of-cond-exp-diff-eq-zero:
assumes adapted: adapted-process M F t₀ X
  and integrable: ∀i. t₀ ≤ i ⇒ integrable M (X i)
  and diff-zero: ∀i j. t₀ ≤ i ≤ j ⇒ AE x in M. cond-exp M (F i) (λξ. X j ξ - X i ξ) x = 0
    shows martingale M F t₀ X
proof
  interpret adapted-process M F t₀ X by (rule adapted)
  {
    fix i j :: 'b assume asm: t₀ ≤ i i ≤ j
    thus AE ξ in M. X i ξ = cond-exp M (F i) (X j) ξ
      using diff-zero[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1),
      of F i j i]
        sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
      fastforce
    }
  qed (auto intro: integrable adapted[THEN adapted-process.adapted])
}

lemma (in sigma-finite-filtered-measure) martingale-of-set-integral-eq:
assumes adapted: adapted-process M F t₀ X
  and integrable: ∀i. t₀ ≤ i ⇒ integrable M (X i)
  and ∀A i j. t₀ ≤ i ≤ j ⇒ A ∈ F i ⇒ set-lebesgue-integral M A (X i) = set-lebesgue-integral M A (X j)
    shows martingale M F t₀ X
proof (unfold-locales)
  fix i j :: 'b assume asm: t₀ ≤ i i ≤ j
  interpret adapted-process M F t₀ X by (rule adapted)
  interpret sigma-finite-subalgebra M F i using asm by blast
  interpret r: sigma-finite-measure restr-to-subalg M (F i) by (simp add: sigma-fin-subalg)
  {
    fix A assume A ∈ restr-to-subalg M (F i)
    hence *: A ∈ F i using sets-restr-to-subalg subalgebras asm by blast
    have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
      M A (X i) using * subalg asm by (auto simp: set-lebesgue-integral-def intro: integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
    also have ... = set-lebesgue-integral M A (cond-exp M (F i) (X j)) using *
      assms(3)[OF asm] cond-exp-set-integral[OF integrable] asm by auto
    finally have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
      (restr-to-subalg M (F i)) A (cond-exp M (F i) (X j)) using * subalg by (auto simp:
      set-lebesgue-integral-def intro!: integral-subalgebra2[symmetric] borel-measurable-scaleR)
  }

```

```

borel-measurable-cond-exp borel-measurable-indicator)
}
hence AE  $\xi$  in restr-to-subalg  $M$  ( $F i$ ).  $X i \xi = cond-exp M (F i) (X j) \xi$  using
asm by (intro r.density-unique-banach, auto intro: integrable-in-subalg subalg
borel-measurable-cond-exp integrable)
thus AE  $\xi$  in  $M$ .  $X i \xi = cond-exp M (F i) (X j) \xi$  using AE-restr-to-subalg[ $OF$ 
subalg] by blast
qed (auto intro: integrable adapted[THEN adapted-process.adapted])

```

## 5.5 Submartingale Lemmas

```

context submartingale
begin

```

```

lemma cond-exp-diff-nonneg:
assumes  $t_0 \leq i \leq j$ 
shows AE  $x$  in  $M$ . cond-exp  $M (F i) (\lambda \xi. X j \xi - X i \xi)$   $x \geq 0$ 
using submartingale-property[ $OF$  assms] assms sigma-finite-subalgebra.cond-exp-diff[ $OF$ 
-integrable(1,1), of -  $j$ ] sigma-finite-subalgebra.cond-exp-F-meas[ $OF$  - integrable
adapted, of  $i$ ] by fastforce

lemma add[intro]:
assumes submartingale  $M F t_0 Y$ 
shows submartingale  $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$ 
proof -
interpret  $Y$ : submartingale  $M F t_0 Y$  by (rule assms)
{
fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
hence AE  $\xi$  in  $M$ .  $X i \xi + Y i \xi \leq cond-exp M (F i) (\lambda x. X j x + Y j x) \xi$ 
using sigma-finite-subalgebra.cond-exp-add[ $OF$  - integrable submartingale.integrable[ $OF$ 
assms], of  $F i j j$ ] submartingale-property[ $OF$  asm] submartingale.submartingale-property[ $OF$ 
assms asm] add-mono[of  $X i$  - -  $Y i$ ] by force
}
thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-add
random-variable adapted integrable Y.random-variable Y.adapted submartingale.integrable)

```

```

qed

```

```

lemma diff[intro]:
assumes supermartingale  $M F t_0 Y$ 
shows submartingale  $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$ 
proof -
interpret  $Y$ : supermartingale  $M F t_0 Y$  by (rule assms)
{
fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
hence AE  $\xi$  in  $M$ .  $X i \xi - Y i \xi \leq cond-exp M (F i) (\lambda x. X j x - Y j x) \xi$ 
using sigma-finite-subalgebra.cond-exp-diff[ $OF$  - integrable supermartingale.integrable[ $OF$ 
assms], of  $F i j j$ ]

```

```

    submartingale-property[OF asm] supermartingale.supermartingale-property[OF
assms asm] diff-mono[of X i - - - Y i - ] by force
}
thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-diff
random-variable adapted integrable Y.random-variable Y.adapted supermartingale.integrable)

```

qed

**lemma** scaleR-nonneg:

assumes  $c \geq 0$

shows submartingale M F t<sub>0</sub> ( $\lambda i \xi. c *_R X i \xi$ )

proof

{

fix i j :: 'b assume asm:  $t_0 \leq i \leq j$

thus AE  $\xi$  in M.  $c *_R X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$

using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i  
j c] submartingale-property[OF asm] by (fastforce intro!: scaleR-left-mono[OF -  
assms])

}

qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable  
random-variable adapted borel-measurable-const-scaleR)

**lemma** scaleR-le-zero:

assumes  $c \leq 0$

shows supermartingale M F t<sub>0</sub> ( $\lambda i \xi. c *_R X i \xi$ )

proof

{

fix i j :: 'b assume asm:  $t_0 \leq i \leq j$

thus AE  $\xi$  in M.  $c *_R X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$

using sigma-finite-subalgebra.cond-exp-scaleR-right[OF - integrable, of F i j  
c] submartingale-property[OF asm]

by (fastforce intro!: scaleR-left-mono-neg[OF - assms])

}

qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable  
random-variable adapted borel-measurable-const-scaleR)

**lemma** uminus[intro]:

shows supermartingale M F t<sub>0</sub> (- X)

unfolding fun-Compl-def using scaleR-le-zero[of -1] by simp

end

**context** submartingale-linorder

begin

**lemma** set-integral-le:

assumes  $A \in F i$   $t_0 \leq i \leq j$

shows set-lebesgue-integral M A (X i)  $\leq$  set-lebesgue-integral M A (X j)

using submartingale-property[OF assms(2), of j] assms subsetD[OF sets-F-subset]

```

by (subst sigma-finite-subalgebra.cond-exp-set-integral[OF - integrable assms(1),
of j])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator
integrable simp add: set-lebesgue-integral-def)

lemma max:
  assumes submartingale M F t₀ Y
  shows submartingale M F t₀ (λi ξ. max (X i ξ) (Y i ξ))
proof (unfold-locales)
  interpret Y: submartingale-linorder M F t₀ Y by (intro submartingale-linorder.intro
assms)
  {
    fix i j :: 'b assume asm: t₀ ≤ i i ≤ j
    have AE ξ in M. max (X i ξ) (Y i ξ) ≤ max (cond-exp M (F i) (X j) ξ)
      (cond-exp M (F i) (Y j) ξ) using submartingale-property Y.submartingale-property
      asm unfolding max-def by fastforce
    thus AE ξ in M. max (X i ξ) (Y i ξ) ≤ cond-exp M (F i) (λξ. max (X j ξ) (Y
j ξ)) ξ using sigma-finite-subalgebra.cond-exp-max[OF - integrable Y.integrable, of
F i j j] asm by (fast intro: order.trans)
  }
  show ∀i. t₀ ≤ i ⇒ (λξ. max (X i ξ) (Y i ξ)) ∈ borel-measurable (F i) ∀i.
t₀ ≤ i ⇒ integrable M (λξ. max (X i ξ) (Y i ξ)) by (force intro: Y.integrable
integrable assms) +
qed

lemma max-0:
  shows submartingale M F t₀ (λi ξ. max 0 (X i ξ))
proof -
  interpret zero: martingale-linorder M F t₀ λ- -. 0 by (force intro: martingale-linorder.intro
martingale-order.intro)
  show ?thesis by (intro zero.max submartingale-linorder.intro submartingale-axioms)
qed

end

lemma (in sigma-finite-filtered-measure) submartingale-of-cond-exp-diff-nonneg:
  assumes adapted: adapted-process M F t₀ X
  and integrable: ∀i. t₀ ≤ i ⇒ integrable M (X i)
  and diff-nonneg: ∀i j. t₀ ≤ i ⇒ i ≤ j ⇒ AE x in M. cond-exp M (F i)
    (λξ. X j ξ - X i ξ) x ≥ 0
  shows submartingale M F t₀ X
proof (unfold-locales)
  interpret adapted-process M F t₀ X by (rule adapted)
  {
    fix i j :: 'b assume asm: t₀ ≤ i i ≤ j
    thus AE ξ in M. X i ξ ≤ cond-exp M (F i) (X j) ξ
      using diff-nonneg[OF asm] sigma-finite-subalgebra.cond-exp-diff[OF - integrable(1,1), of F i j i]
        sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable adapted, of i] by
  }

```

```

fastforce
}
qed (auto intro: integrable adapted[THEN adapted-process.adapted])

lemma (in sigma-finite-filtered-measure) submartingale-of-set-integral-le:
  fixes X :: - ⇒ - ⇒ - :: {linorder-topology}
  assumes adapted: adapted-process M F t₀ X
    and integrable: ∀ i. t₀ ≤ i ⇒ integrable M (X i)
    and ∀ A i j. t₀ ≤ i ⇒ i ≤ j ⇒ A ∈ F i ⇒ set-lebesgue-integral M A (X
i) ≤ set-lebesgue-integral M A (X j)
    shows submartingale M F t₀ X
proof (unfold-locales)
{
  fix i j :: 'b assume asm: t₀ ≤ i i ≤ j
  interpret adapted-process M F t₀ X by (rule adapted)
  interpret r: sigma-finite-measure restr-to-subalg M (F i) using asm sigma-finite-subalgebra.sigma-fin-subalg
by blast
{
  fix A assume A ∈ restr-to-subalg M (F i)
  hence *: A ∈ F i using asm sets-restr-to-subalg subalgebras by blast
  have set-lebesgue-integral (restr-to-subalg M (F i)) A (X i) = set-lebesgue-integral
M A (X i) using * asm subalgebras by (auto simp: set-lebesgue-integral-def intro:
integral-subalgebra2 borel-measurable-scaleR adapted borel-measurable-indicator)
  also have ... ≤ set-lebesgue-integral M A (cond-exp M (F i) (X j)) using
* assms(3)[OF asm] asm sigma-finite-subalgebra.cond-exp-set-integral[OF - inte-
grable] by fastforce
  also have ... = set-lebesgue-integral (restr-to-subalg M (F i)) A (cond-exp M
(F i) (X j)) using * asm subalgebras by (auto simp: set-lebesgue-integral-def intro!:
integral-subalgebra2[symmetric] borel-measurable-scaleR borel-measurable-cond-exp
borel-measurable-indicator)
  finally have 0 ≤ set-lebesgue-integral (restr-to-subalg M (F i)) A (λξ. cond-exp
M (F i) (X j) ξ − X i ξ) using * asm subalgebras by (subst set-integral-diff,
auto simp add: set-integrable-def sets-restr-to-subalg intro!: integrable adapted inte-
grable-in-subalg borel-measurable-scaleR borel-measurable-indicator borel-measurable-cond-exp
integrable-mult-indicator)
}
hence AE ξ in restr-to-subalg M (F i). 0 ≤ cond-exp M (F i) (X j) ξ − X i ξ
  by (intro r.density-nonneg integrable-in-subalg asm subalgebras borel-measurable-diff
borel-measurable-cond-exp adapted Bochner-Integration.integrable-diff integrable-cond-exp
integrable)
  thus AE ξ in M. X i ξ ≤ cond-exp M (F i) (X j) ξ using AE-restr-to-subalg[OF
subalgebras] asm by simp
}
qed (auto intro: integrable adapted[THEN adapted-process.adapted])

```

## 5.6 Supermartingale Lemmas

The following lemmas are exact duals of the ones for submartingales.

**context** supermartingale

```

begin

lemma cond-exp-diff-nonneg:
assumes  $t_0 \leq i \leq j$ 
shows  $\text{AE } x \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X i \xi - X j \xi) x \geq 0$ 
using assms supermartingale-property[ $\text{OF assms}$ ] sigma-finite-subalgebra.cond-exp-diff[ $\text{OF}$ -integrable(1,1), of  $F i i j$ ]
sigma-finite-subalgebra.cond-exp-F-meas[ $\text{OF}$ -integrable adapted, of  $i$ ] by
fastforce

lemma add[intro]:
assumes supermartingale  $M F t_0 Y$ 
shows supermartingale  $M F t_0 (\lambda i \xi. X i \xi + Y i \xi)$ 
proof -
interpret  $Y$ : supermartingale  $M F t_0 Y$  by (rule assms)
{
fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
hence  $\text{AE } \xi \text{ in } M. X i \xi + Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x + Y j x) \xi$ 
using sigma-finite-subalgebra.cond-exp-add[ $\text{OF}$ -integrable supermartingale.integrable[ $\text{OF assms}$ ], of  $F i j j$ ]
supermartingale-property[ $\text{OF asm}$ ] supermartingale.supermartingale-property[ $\text{OF assms asm}$ ] add-mono[of -  $X i - Y i$ ] by force
}
thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-add
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted supermartingale.integrable)

qed

lemma diff[intro]:
assumes submartingale  $M F t_0 Y$ 
shows supermartingale  $M F t_0 (\lambda i \xi. X i \xi - Y i \xi)$ 
proof -
interpret  $Y$ : submartingale  $M F t_0 Y$  by (rule assms)
{
fix  $i j :: 'b$  assume asm:  $t_0 \leq i \leq j$ 
hence  $\text{AE } \xi \text{ in } M. X i \xi - Y i \xi \geq \text{cond-exp } M (F i) (\lambda x. X j x - Y j x) \xi$ 
using sigma-finite-subalgebra.cond-exp-diff[ $\text{OF}$ -integrable submartingale.integrable[ $\text{OF assms}$ ], of  $F i j j$ , unfolded fun-diff-def]
supermartingale-property[ $\text{OF asm}$ ] submartingale.supermartingale-property[ $\text{OF assms asm}$ ] diff-mono[of -  $X i - Y i$ ] by force
}
thus ?thesis using assms by (unfold-locales) (auto simp add: borel-measurable-diff
random-variable adapted integrable  $Y$ .random-variable  $Y$ .adapted submartingale.integrable)

qed

lemma scaleR-nonneg:
assumes  $c \geq 0$ 
shows supermartingale  $M F t_0 (\lambda i \xi. c *_R X i \xi)$ 

```

```

proof
{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  thus  $\text{AE } \xi \text{ in } M. c *_R X i \xi \geq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
    using sigma-finite-subalgebra.cond-exp-scaleR-right[ $\text{OF - integrable, of } F i j c$ ] supermartingale-property[ $\text{OF asm}$ ] by (fastforce intro!: scaleR-left-mono[ $\text{OF - assms}$ ])
}
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR)

lemma scaleR-le-zero:
assumes  $c \leq 0$ 
shows submartingale  $M F t_0 (\lambda i \xi. c *_R X i \xi)$ 
proof
{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  thus  $\text{AE } \xi \text{ in } M. c *_R X i \xi \leq \text{cond-exp } M (F i) (\lambda \xi. c *_R X j \xi) \xi$ 
    using sigma-finite-subalgebra.cond-exp-scaleR-right[ $\text{OF - integrable, of } F i j c$ ] supermartingale-property[ $\text{OF asm}$ ] by (fastforce intro!: scaleR-left-mono-neg[ $\text{OF - assms}$ ])
}
qed (auto simp add: borel-measurable-integrable borel-measurable-scaleR integrable random-variable adapted borel-measurable-const-scaleR)

lemma uminus[intro]:
shows submartingale  $M F t_0 (- X)$ 
unfolding fun-Compl-def using scaleR-le-zero[of  $-1$ ] by simp

end

context supermartingale-linorder
begin

lemma set-integral-ge:
assumes  $A \in F i$   $t_0 \leq i \leq j$ 
shows set-lebesgue-integral  $M A (X i) \geq \text{set-lebesgue-integral } M A (X j)$ 
using supermartingale-property[ $\text{OF assms}(2)$ , of  $j$ ] assms subsetD[ $\text{OF sets-F-subset}$ ]
by (subst sigma-finite-subalgebra.cond-exp-set-integral[ $\text{OF - integrable assms}(1)$ , of  $j$ ])
  (auto intro!: scaleR-left-mono integral-mono-AE-banach integrable-mult-indicator
integrable simp add: set-lebesgue-integral-def)

lemma min:
assumes supermartingale  $M F t_0 Y$ 
shows supermartingale  $M F t_0 (\lambda i \xi. \min (X i \xi) (Y i \xi))$ 
proof (unfold-locales)
  interpret  $Y$ : supermartingale-linorder  $M F t_0 Y$  by (intro supermartingale-linorder.intro assms)

```

```

{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  have  $\text{AE } \xi \text{ in } M. \min(X_i \xi) (Y_i \xi) \geq \min(\text{cond-exp } M(F_i) (X_j) \xi) (\text{cond-exp } M(F_i) (Y_j) \xi)$  using supermartingale-property  $Y.\text{supermartingale-property}$   $\text{asm}$ 
  unfolding min-def by fastforce
  thus  $\text{AE } \xi \text{ in } M. \min(X_i \xi) (Y_i \xi) \geq \text{cond-exp } M(F_i) (\lambda \xi. \min(X_j \xi) (Y_j \xi)) \xi$  using sigma-finite-subalgebra.cond-exp-min[ $OF$  - integrable  $Y.\text{integrable}$ , of  $F_i j j$ ]  $\text{asm}$  by (fast intro: order.trans)
}
show  $\bigwedge i. t_0 \leq i \implies (\lambda \xi. \min(X_i \xi) (Y_i \xi)) \in \text{borel-measurable}(F_i) \bigwedge i. t_0 \leq i \implies \text{integrable } M(\lambda \xi. \min(X_i \xi) (Y_i \xi))$  by (force intro:  $Y.\text{integrable}$  integrable assms) +
qed

lemma min-0:
  shows supermartingale  $M F t_0 (\lambda i \xi. \min 0 (X_i \xi))$ 
proof -
  interpret zero: martingale-linorder  $M F t_0 \lambda \cdot \cdot 0$  by (force intro: martingale-linorder.intro)
  show ?thesis by (intro zero.min supermartingale-linorder.intro supermartingale-axioms)
qed

end

lemma (in sigma-finite-filtered-measure) supermartingale-of-cond-exp-diff-le-zero:
  assumes adapted: adapted-process  $M F t_0 X$ 
  and integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M(X_i)$ 
  and diff-le-zero:  $\bigwedge i j. t_0 \leq i \implies i \leq j \implies \text{AE } x \text{ in } M. \text{cond-exp } M(F_i) (\lambda \xi. X_j \xi - X_i \xi) x \leq 0$ 
  shows supermartingale  $M F t_0 X$ 
proof
  interpret adapted-process  $M F t_0 X$  by (rule adapted)
{
  fix i j :: 'b assume asm:  $t_0 \leq i \leq j$ 
  thus  $\text{AE } \xi \text{ in } M. X_i \xi \geq \text{cond-exp } M(F_i) (X_j) \xi$ 
  using diff-le-zero[ $OF$   $\text{asm}$ ] sigma-finite-subalgebra.cond-exp-diff[ $OF$  - integrable(1,1), of  $F_i j i$ ]
    sigma-finite-subalgebra.cond-exp-F-meas[ $OF$  - integrable adapted, of  $i$ ] by
  fastforce
}
qed (auto intro: integrable adapted[THEN adapted-process.adapted])

lemma (in sigma-finite-filtered-measure) supermartingale-of-set-integral-ge:
  fixes  $X :: - \Rightarrow - \Rightarrow - :: \{\text{linorder-topology}\}$ 
  assumes adapted: adapted-process  $M F t_0 X$ 
  and integrable:  $\bigwedge i. t_0 \leq i \implies \text{integrable } M(X_i)$ 
  and  $\bigwedge A i j. t_0 \leq i \leq j \implies A \in F_i \implies \text{set-lebesgue-integral } M A (X_j) \leq \text{set-lebesgue-integral } M A (X_i)$ 

```

```

shows supermartingale M F t0 X
proof –
  interpret adapted-process M F t0 X by (rule adapted)
  note * = set-integral-uminus[unfolded set-integrable-def, OF integrable-mult-indicator[OF - integrable]]
  have supermartingale M F t0 (-( - X))
  using ord-eq-le-trans[OF * ord-le-eq-trans[OF le-imp-neg-le[OF assms(3)] *[symmetric]]]
  sets-F-subset[THEN subsetD]
  by (intro submartingale.uminus submartingale-of-set-integral-le[OF uminus-adapted])

  (clarsimp simp add: fun-Compl-def integrable | fastforce) +
  thus ?thesis unfolding fun-Compl-def by simp
qed

```

Many of the statements we have made concerning martingales can be simplified when the indexing set is the natural numbers. Given a point in time  $i \in \mathbb{N}$ , it suffices to consider the successor  $i + 1$ , instead of all future times  $i \leq j$ .

## 5.7 Discrete Time Martingales

```

context nat-sigma-finite-filtered-measure
begin

```

A predictable martingale is necessarily constant.

```

lemma predictable-const:
  assumes martingale M F 0 X
  and predictable-process M F 0 X
  shows AE  $\xi$  in M.  $X i \xi = X j \xi$ 
proof –
  interpret martingale M F 0 X by (rule assms)
  have *: AE  $\xi$  in M.  $X i \xi = X 0 \xi$  for i
  proof (induction i)
    case 0
    then show ?case by (simp add: bot-nat-def)
  next
    case (Suc i)
    interpret S: adapted-process M F 0  $\lambda i. X (Suc i)$  by (intro predictable-implies-adapted-Suc assms)
    show ?case using Suc S.adapted[of i] martingale-property[OF - le-SucI, of i]
    sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc i] by fastforce
    qed
    show ?thesis using *[of i] *[of j] by force
  qed

lemma martingale-of-set-integral-eq-Suc:
  assumes adapted: adapted-process M F 0 X
  and integrable:  $\bigwedge i. \text{integrable } M (X i)$ 

```

```

and  $\bigwedge A \ i. A \in F \ i \implies \text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ (\text{Suc } i))$ 
shows martingale  $M \ F \ 0 \ X$ 
proof (intro martingale-of-set-integral-eq adapted integrable)
fix  $i \ j \ A$  assume  $\text{asm}: i \leq j \ A \in \text{sets } (F \ i)$ 
show  $\text{set-lebesgue-integral } M \ A \ (X \ i) = \text{set-lebesgue-integral } M \ A \ (X \ j)$  using
asm
proof (induction  $j - i$  arbitrary:  $i \ j$ )
case 0
then show ?case using asm by simp
next
case ( $\text{Suc } n$ )
hence  $*: n = j - \text{Suc } i$  by linarith
have  $\text{Suc } i \leq j$  using  $\text{Suc}(2,3)$  by linarith
thus ?case using sets-F-mono[ $OF - \text{le-SucI}$ ]  $\text{Suc}(4)$   $\text{Suc}(1)[OF *$ ] by (auto
intro: assms(3)[THEN trans])
qed
qed

lemma martingale-nat:
assumes adapted: adapted-process  $M \ F \ 0 \ X$ 
and integrable:  $\bigwedge i. \text{integrable } M \ (X \ i)$ 
and  $\bigwedge i. AE \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (\text{Suc } i)) \ \xi$ 
shows martingale  $M \ F \ 0 \ X$ 
proof (unfold-locales)
interpret adapted-process  $M \ F \ 0 \ X$  by (rule adapted)
fix  $i \ j :: nat$  assume  $\text{asm}: i \leq j$ 
show  $AE \xi \text{ in } M. X \ i \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi$  using asm
proof (induction  $j - i$  arbitrary:  $i \ j$ )
case 0
hence  $j = i$  by simp
thus ?case using sigma-finite-subalgebra.cond-exp-F-meas[ $OF - \text{integrable adapted}$ ,
THEN  $AE\text{-symmetric}$ ] by blast
next
case ( $\text{Suc } n$ )
have  $j: j = \text{Suc } (n + i)$  using Suc by linarith
have  $n: n = n + i - i$  using Suc by linarith
have  $*: AE \xi \text{ in } M. \text{cond-exp } M \ (F \ (n + i)) \ (X \ j) \ \xi = X \ (n + i) \ \xi$  unfolding
j using assms(3)[THEN  $AE\text{-symmetric}$ ] by blast
have  $AE \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi = \text{cond-exp } M \ (F \ i) \ (\text{cond-exp } M
(F \ (n + i)) \ (X \ j)) \ \xi$  by (intro cond-exp-nested-subalg integrable subalg, simp add:
subalgebra-def sets-F-mono)
hence  $AE \xi \text{ in } M. \text{cond-exp } M \ (F \ i) \ (X \ j) \ \xi = \text{cond-exp } M \ (F \ i) \ (X \ (n + i))$ 
\xi using cond-exp-cong-AE[ $OF \text{ integrable-cond-exp integrable } *$ ] by force
thus ?case using  $\text{Suc}(1)[OF \ n]$  by fastforce
qed
qed (auto simp add: integrable adapted[THEN adapted-process.adapted])

lemma martingale-of-cond-exp-diff-Suc-eq-zero:

```

```

assumes adapted: adapted-process M F 0 X
  and integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and  $\bigwedge i. AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X (\text{Suc } i) \xi - X i \xi) \xi = 0$ 
  shows martingale M F 0 X
proof (intro martingale-nat integrable adapted)
  interpret adapted-process M F 0 X by (rule adapted)
  fix i
  show AE  $\xi$  in M.  $X i \xi = \text{cond-exp } M (F i) (X (\text{Suc } i)) \xi$  using cond-exp-diff[OF
integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(3)[of
i] by fastforce
qed

end

```

## 5.8 Discrete Time Submartingales

```

context nat-sigma-finite-filtered-measure
begin

lemma predictable-mono:
assumes submartingale M F 0 X
  and predictable-process M F 0 X  $i \leq j$ 
  shows AE  $\xi$  in M.  $X i \xi \leq X j \xi$ 
  using assms(3)
proof (induction j - i arbitrary: i j)
  case 0
  then show ?case by simp
next
  case (Suc n)
  hence *:  $n = j - \text{Suc } i$  by linarith
  interpret submartingale M F 0 X by (rule assms)
  interpret S: adapted-process M F 0  $\lambda i. X (\text{Suc } i)$  by (intro predictable-implies-adapted-Suc
assms)
  have Suc i  $\leq j$  using Suc(2,3) by linarith
  thus ?case using Suc(1)[OF *] S.adapted[of i] submartingale-property[OF -
le-SucI, of i] sigma-finite-subalgebra.cond-exp-F-meas[OF - integrable, of F i Suc
i] by fastforce
qed

lemma submartingale-of-set-integral-le-Suc:
fixes X :: -  $\Rightarrow$  -  $\Rightarrow$  - :: {linorder-topology}
assumes adapted: adapted-process M F 0 X
  and integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and  $\bigwedge A i. A \in F i \implies \text{set-lebesgue-integral } M A (X i) \leq \text{set-lebesgue-integral }$ 
 $M A (X (\text{Suc } i))$ 
  shows submartingale M F 0 X
proof (intro submartingale-of-set-integral-le adapted integrable)
  fix i j A assume asm:  $i \leq j A \in \text{sets } (F i)$ 
  show set-lebesgue-integral M A (X i)  $\leq \text{set-lebesgue-integral } M A (X j)$  using

```

```


asm
  proof (induction j - i arbitrary: i j)
    case 0
    then show ?case using asm by simp
  next
    case (Suc n)
    hence *: n = j - Suc i by linarith
    have Suc i ≤ j using Suc(2,3) by linarith
    thus ?case using sets-F-mono[OF - le-SucI] Suc(4) Suc(1)[OF *] by (auto
      intro: assms(3)[THEN order-trans])
    qed
  qed

lemma submartingale-nat:
  fixes X :: - ⇒ - ⇒ - :: {linorder-topology}
  assumes adapted: adapted-process M F 0 X
    and integrable: ∀i. integrable M (X i)
    and ∀i. AE ξ in M. X i ξ ≤ cond-exp M (F i) (X (Suc i)) ξ
    shows submartingale M F 0 X
proof -
  show ?thesis using subalg assms(3) integrable
    by (intro submartingale-of-set-integral-le-Suc adapted integrable ord-le-eq-trans[OF
      set-integral-mono-AE-banach cond-exp-set-integral[symmetric]])
      (meson in-mono integrable-mult-indicator set-integrable-def subalgebra-def,
       meson integrable-cond-exp in-mono integrable-mult-indicator set-integrable-def sub-
       algebra-def, fast+)
  qed

lemma submartingale-of-cond-exp-diff-Suc-nonneg:
  fixes X :: - ⇒ - ⇒ - :: {linorder-topology}
  assumes adapted: adapted-process M F 0 X
    and integrable: ∀i. integrable M (X i)
    and ∀i. AE ξ in M. cond-exp M (F i) (λξ. X (Suc i)) ξ - X i ξ ≥ 0
    shows submartingale M F 0 X
proof (intro submartingale-nat integrable adapted)
  interpret adapted-process M F 0 X by (rule assms)
  fix i
  show AE ξ in M. X i ξ ≤ cond-exp M (F i) (X (Suc i)) ξ using cond-exp-diff[OF
    integrable(1,1), of i Suc i i] cond-exp-F-meas[OF integrable adapted, of i] assms(3)[of
    i] by fastforce
  qed

lemma submartingale-partial-sum-scaleR:
  assumes submartingale-linorder M F 0 X
    and adapted-process M F 0 C ∀i. AE ξ in M. 0 ≤ C i ξ ∧i. AE ξ in M. C i
    ξ ≤ R
    shows submartingale M F 0 (λn ξ. ∑ i< n. C i ξ *R (X (Suc i)) ξ - X i ξ))
proof -
  interpret submartingale-linorder M F 0 X by (rule assms)


```

```

interpret C: adapted-process M F 0 C by (rule assms)
interpret C': adapted-process M F 0 λi. C (i - 1) ξ *R (X i ξ - X (i - 1) ξ) by (intro adapted-process.scaleR-right-adapted adapted-process.diff-adapted, unfold-locales) (auto intro: adaptedD C.adaptedD)+
interpret S: adapted-process M F 0 λn. Σ i < n. C i ξ *R (X (Suc i) ξ - X i ξ) using C'.adapted-process-axioms[THEN partial-sum-Suc-adapted] diff-Suc-1 by simp
have integrable M (λx. C i x *R (X (Suc i) x - X i x)) for i using assms(3,4)[of i] by (intro Bochner-Integration.integrable-bound[OF integrable-scaleR-right, OF Bochner-Integration.integrable-diff, OF integrable(1,1), of R Suc i i]) (auto simp add: mult-mono)
moreover have AE ξ in M. 0 ≤ cond-exp M (F i) (λξ. (Σ i < Suc i. C i ξ *R (X (Suc i) ξ - X i ξ)) - (Σ i < i. C i ξ *R (X (Suc i) ξ - X i ξ))) ξ for i
  using sigma-finite-subalgebra.cond-exp-measurable-scaleR[OF - calculation - C.adapted, of i]
  cond-exp-diff-nonneg[OF - le-SucI, OF - order.refl, of i] assms(3,4)[of i]
by (fastforce simp add: scaleR-nonneg-nonneg integrable)
ultimately show ?thesis by (intro submartingale-of-cond-exp-diff-Suc-nonneg S.adapted-process-axioms Bochner-Integration.integrable-sum, blast+)
qed

lemma submartingale-partial-sum-scaleR':
assumes submartingale-linorder M F 0 X
  and predictable-process M F 0 C ∧i. AE ξ in M. 0 ≤ C i ξ ∧i. AE ξ in M.
C i ξ ≤ R
  shows submartingale M F 0 (λn. Σ i < n. C (Suc i) ξ *R (X (Suc i) ξ - X i ξ))
proof -
interpret Suc-C: adapted-process M F 0 λi. C (Suc i) using predictable-implies-adapted-Suc assms by blast
show ?thesis by (intro submartingale-partial-sum-scaleR[OF assms(1), of - R] assms) (intro-locales)
qed

end

```

## 5.9 Discrete Time Supermartingales

```

context nat-sigma-finite-filtered-measure
begin

```

```

lemma predictable-mono':
assumes supermartingale M F 0 X
  and predictable-process M F 0 X i ≤ j
  shows AE ξ in M. X i ξ ≥ X j ξ
  using assms(3)
proof (induction j - i arbitrary: i j)
  case 0
  then show ?case by simp

```

```

next
  case ( $Suc n$ )
    hence  $*: n = j - Suc i$  by linarith
    interpret supermartingale  $M F 0 X$  by (rule assms)
    interpret  $S: adapted\text{-process}$   $M F 0 \lambda i. X (Suc i)$  by (intro predictable-implies-adapted-Suc assms)
    have  $Suc i \leq j$  using  $Suc(2,3)$  by linarith
    thus ?case using  $Suc(1)[OF *] S.adapted[of i]$  supermartingale-property[ $OF - le\text{-}SucI$ , of  $i$ ] sigma-finite-subalgebra.cond-exp-F-meas[ $OF - integrable$ , of  $F i Suc i$ ] by fastforce
  qed

lemma supermartingale-of-set-integral-ge-Suc:
  fixes  $X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$ 
  assumes adapted: adapted-process  $M F 0 X$ 
    and integrable:  $\bigwedge i. integrable M (X i)$ 
    and  $\bigwedge A i. A \in F i \implies set\text{-}lebesgue\text{-}integral M A (X i) \geq set\text{-}lebesgue\text{-}integral M A (X (Suc i))$ 
    shows supermartingale  $M F 0 X$ 
proof –
  interpret adapted-process  $M F 0 X$  by (rule assms)
  interpret uminus-X: adapted-process  $M F 0 -X$  by (rule uminus-adapted)
  note  $* = set\text{-}integral\text{-}uminus[unfolded set\text{-}integrable\text{-}def, OF integrable\text{-}mult\text{-}indicator[ $OF - integrable$ ]]]$ 
  have supermartingale  $M F 0 (-(- X))$ 
  using ord-eq-le-trans[ $OF * ord\text{-}le\text{-}eq\text{-}trans[OF le\text{-}imp\text{-}neg\text{-}le[ $OF assms(3)$ ]] *[symmetric]]]
  sets-F-subset[THEN subsetD]
    by (intro submartingale.uminus submartingale-of-set-integral-le-Suc[ $OF uminus\text{-}adapted$ ])
    (clarsimp simp add: fun-Compl-def integrable | fastforce)+
  thus ?thesis unfolding fun-Compl-def by simp
  qed

lemma supermartingale-nat:
  fixes  $X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$ 
  assumes adapted: adapted-process  $M F 0 X$ 
    and integrable:  $\bigwedge i. integrable M (X i)$ 
    and  $\bigwedge i. AE \xi \text{ in } M. X i \xi \geq cond\text{-}exp M (F i) (X (Suc i)) \xi$ 
    shows supermartingale  $M F 0 X$ 
proof –
  interpret adapted-process  $M F 0 X$  by (rule assms)
  have  $AE \xi \text{ in } M. -X i \xi \leq cond\text{-}exp M (F i) (\lambda x. -X (Suc i) x) \xi$  for  $i$  using assms(3) cond-exp-uminus[ $OF integrable$ , of  $i Suc i$ ] by force
  hence supermartingale  $M F 0 (-(- X))$  by (intro submartingale.uminus submartingale-nat[ $OF uminus\text{-}adapted$ ]) (auto simp add: fun-Compl-def integrable)
  thus ?thesis unfolding fun-Compl-def by simp
  qed

lemma supermartingale-of-cond-exp-diff-Suc-le-zero:$ 
```

```

fixes  $X :: - \Rightarrow - \Rightarrow - :: \{linorder-topology\}$ 
assumes adapted: adapted-process  $M F 0 X$ 
  and integrable:  $\bigwedge i. \text{integrable } M (X i)$ 
  and  $\bigwedge i. AE \xi \text{ in } M. \text{cond-exp } M (F i) (\lambda \xi. X (\text{Suc } i) \xi - X i \xi) \xi \leq 0$ 
  shows supermartingale  $M F 0 X$ 
proof (intro supermartingale-nat integrable adapted)
interpret adapted-process  $M F 0 X$  by (rule assms)
fix  $i$ 
show  $AE \xi \text{ in } M. X i \xi \geq \text{cond-exp } M (F i) (X (\text{Suc } i)) \xi$  using cond-exp-diff[ $OF$ 
integrable(1,1), of  $i$  Suc  $i$   $i$ ] cond-exp-F-meas[ $OF$  integrable adapted, of  $i$ ] assms(3)[of
 $i$ ] by fastforce
qed

end

end

```

```

theory Example-Coin-Toss
imports Martingale HOL-Probability.Stream-Space HOL-Probability.Probability-Mass-Function
begin

```

## 6 Example: Coin Toss

Consider a coin-tossing game, where the coin lands on heads with probability  $p \in [0, 1]$ . Assume that the gambler wins a fixed amount  $c > 0$  on a heads outcome and loses the same amount  $c$  on a tails outcome. Let  $(X_n)_{n \in \mathbb{N}}$  be a stochastic process, where  $X_n$  denotes the gambler's fortune after the  $n$ -th coin toss. Then, we have the following three cases.

1. If  $p = 1/2$ , it means the coin is fair and has an equal chance of landing heads or tails. In this case, the gambler, on average, neither wins nor loses money over time. The expected value of the gambler's fortune stays the same over time. Therefore,  $(X_n)_{n \in \mathbb{N}}$  is a martingale.
2. If  $p \geq 1/2$ , it means the coin is biased in favor of heads. In this case, the gambler is more likely to win money on each bet. Over time, the gambler's fortune tends to increase on average. Therefore,  $(X_n)_{n \in \mathbb{N}}$  is a submartingale.
3. If  $p \leq 1/2$ , it means the coin is biased in favor of tails. In this scenario, the gambler is more likely to lose money on each bet. Over time, the gambler's fortune decreases on average. Therefore,  $(X_n)_{n \in \mathbb{N}}$  is a supermartingale.

To formalize this example, we first consider a probability space consisting of infinite sequences of coin tosses.

```

definition bernoulli-stream :: real  $\Rightarrow$  (bool stream) measure where
  bernoulli-stream p = stream-space (measure-pmf (bernoulli-pmf p))

lemma space-bernoulli-stream[simp]: space (bernoulli-stream p) = UNIV by (simp
add: bernoulli-stream-def space-stream-space)

We define the fortune of the player at time n to be the number of heads
minus number of tails.

definition fortune :: nat  $\Rightarrow$  bool stream  $\Rightarrow$  real where
  fortune n = ( $\lambda s. \sum b \leftarrow \text{stake} (\text{Suc } n) s. \text{if } b \text{ then } 1 \text{ else } -1$ )

definition toss :: nat  $\Rightarrow$  bool stream  $\Rightarrow$  real where
  toss n = ( $\lambda s. \text{if } \text{snth } s n \text{ then } 1 \text{ else } -1$ )

lemma toss-indicator-def: toss n = indicator {s. s !! n} – indicator {s.  $\neg s$  !! n}
  unfolding toss-def indicator-def by force

lemma range-toss: range (toss n) = {-1, 1}
proof –
  have sconst True !! n by simp
  moreover have  $\neg$ sconst False !! n by simp
  ultimately have  $\exists x. x$  !! n  $\exists x. \neg x$  !! n by blast+
  thus ?thesis unfolding toss-def image-def by auto
qed

lemma vimage-toss: toss n – ‘ A = (if  $1 \in A$  then {s. s !! n} else {})  $\cup$  (if  $-1 \in A$  then {s.  $\neg s$  !! n} else {})
  unfolding vimage-def toss-def by auto

lemma fortune-Suc: fortune (Suc n) s = fortune n s + toss (Suc n) s
  by (induction n arbitrary: s) (simp add: fortune-def toss-def)+

lemma fortune-toss-sum: fortune n s = ( $\sum i \in \{\dots n\}. \text{toss } i s$ )
  by (induction n arbitrary: s) (simp add: fortune-def toss-def, simp add: fortune-Suc)

lemma fortune-bound: norm (fortune n s)  $\leq$  Suc n by (induction n) (force simp
add: fortune-toss-sum toss-def)+

```

Our definition of *bernoulli-stream* constitutes a probability space.

**interpretation** prob-space bernoulli-stream *p* **unfolding** bernoulli-stream-def **by**
(simp add: measure-pmf.prob-space-axioms prob-space.prob-space-stream-space)

**abbreviation** toss-filtration *p*  $\equiv$  nat-natural-filtration (bernoulli-stream *p*) toss

The stochastic process *toss* is adapted to the filtration it generates.

**interpretation** toss: adapted-process bernoulli-stream *p* toss-filtration *p* 0 toss
 **by** (intro adapted-process.intro stochastic-process.adapted-process-natural-filtration)
 (unfold-locales, auto simp add: toss-def bernoulli-stream-def)

```

interpretation bernoulli-stream-natural-filtration: nat-finite-filtered-measure bernoulli-stream
p toss-filtration p
by (simp add: nat-finite-filtered-measure-def toss.finite-filtered-measure-natural-filtration)

```

Similarly, the stochastic process *fortune* is adapted to the filtration generated by the tosses.

```

interpretation fortune: adapted-process bernoulli-stream p toss-filtration p 0 for-
tune

```

**proof** –

```
show adapted-process (bernoulli-stream p) (toss-filtration p) 0 fortune
```

```
unfoldng fortune-toss-sum
```

```
by (intro toss.partial-sum-adapted[folded atMost-atLeast0]) intro-locales
```

qed

```

lemma integrable-toss: integrable (bernoulli-stream p) (toss n)

```

```
using toss.random-variable
```

```
by (intro Bochner-Integration.integrable-bound[OF integrable-const[of - 1 :: real]]) (auto simp add: toss-def)
```

```

lemma integrable-fortune: integrable (bernoulli-stream p) (fortune n) using for-
tune-bound

```

```
by (intro Bochner-Integration.integrable-bound[OF integrable-const[of - Suc n] fortune.random-variable]) auto
```

We provide the following lemma to explicitly calculate the probability of events in this probability space.

```

lemma measure-bernoulli-stream-snth-pred:

```

```
assumes 0 ≤ p and p ≤ 1 and finite J
```

```
shows prob p {w ∈ space (bernoulli-stream p). ∀ j ∈ J. P j = w !! j} = p ^ card (J ∩ Collect P) * (1 - p) ^ (card (J - Collect P))
```

**proof** –

```
let ?PiE = (Π_E i ∈ J. if P i then {True} else {False})
```

```
have product-prob-space (λ-. measure-pmf (bernoulli-pmf p)) by unfold-locales
```

```

hence *: to-stream -` {s. ∀ i ∈ J. P i = s !! i} = {s. ∀ i ∈ J. P i = s i} using
assms by (simp add: to-stream-def)

```

```
also have ... = prod-emb UNIV (λ-. measure-pmf (bernoulli-pmf p)) J ?PiE
```

**proof** –

```
{
```

```
fix s assume (∀ i ∈ J. P i = s i)
```

```
hence (∀ i ∈ J. P i = s i) = (s ∈ prod-emb UNIV (λ-. measure-pmf (bernoulli-pmf p)) J ?PiE)
```

```
by (subst prod-emb-iff[of s]) (smt (verit, best) not-def assms(3) id-def
PiE-eq-singleton UNIV-I extensional-UNIV insert-iff singletonD space-measure-pmf)
```

```
}
```

```
moreover
```

```
{
```

```
fix s assume ¬(∀ i ∈ J. P i = s i)
```

```

then obtain i where  $i \in J$   $P i \neq s$   $i$  by blast
hence  $(\forall i \in J. P i = s) = (s \in \text{prod-emb UNIV} (\lambda-. \text{measure-pmf} (\text{bernoulli-pmf } p))) J ?PiE)$ 
by (simp add: restrict-def prod-emb-iff[of s]) (smt (verit, ccfv-SIG) PiE-mem
assms(3) id-def insert-iff singleton-iff)
}
ultimately show ?thesis by auto
qed
finally have inteq:  $(\text{to-stream} -^c \{s. \forall i \in J. P i = s !! i\}) = \text{prod-emb UNIV}$ 
 $(\lambda-. \text{measure-pmf} (\text{bernoulli-pmf } p)) J ?PiE$  .
let ?M =  $(\text{Pi}_M \text{ UNIV} (\lambda-. \text{measure-pmf} (\text{bernoulli-pmf } p)))$ 
have emeasure (bernoulli-stream p) { $s \in \text{space} (\text{bernoulli-stream } p)$ .  $\forall i \in J. P i = s !! i$ } = emeasure ?M  $(\text{to-stream} -^c \{s. \forall i \in J. P i = s !! i\})$ 
using assms emeasure-distr[of to-stream ?M (vimage-algebra (streams (space
(measure-pmf (bernoulli-pmf p)))) (!! ?M) {s.  $\forall i \in J. P i = s !! i$ }, symmetric]
measurable-to-stream[of (measure-pmf (bernoulli-pmf p)))]
by (simp only: bernoulli-stream-def stream-space-def *, simp add: space-PiM )
(smt (verit, best) emeasure-notin-sets in-vimage-algebra inf-top.right-neutral sets-distr
vimage-Collect)
also have ... = emeasure ?M  $(\text{prod-emb UNIV} (\lambda-. \text{measure-pmf} (\text{bernoulli-pmf } p))) J ?PiE)$  using inteq by (simp add: space-PiM )
also have ... =  $(\prod_{i \in J} \text{emeasure} (\text{measure-pmf} (\text{bernoulli-pmf } p)))$  (if  $P i$  then
{True} else {False}))
by (subst emeasure-PiM-emb) (auto simp add: prob-space-measure-pmf assms(3))
also have ... =  $(\prod_{i \in J} \text{Collect } P. \text{ennreal } p) * (\prod_{i \in J} \text{Collect } P. \text{ennreal } (1 - p))$ 
unfolding emeasure-pmf-single[of bernoulli-pmf p True, unfolded pmf-bernoulli-True[OF
assms(1,2)], symmetric]
emeasure-pmf-single[of bernoulli-pmf p False, unfolded pmf-bernoulli-False[OF
assms(1,2)], symmetric]
by (simp add: prod.Int-Diff[OF assms(3), of - Collect P])
also have ... =  $p \hat{\wedge} \text{card} (J \cap \text{Collect } P) * (1 - p) \hat{\wedge} \text{card} (J - \text{Collect } P)$  using
assms by (simp add: prod-ennreal ennreal-mult' ennreal-power)
finally show ?thesis using assms by (intro measure-eq-emeasure-eq-ennreal)
auto
qed

```

### lemma

```

assumes  $0 \leq p$  and  $p \leq 1$ 
shows measure-bernoulli-stream-snth: prob p { $w \in \text{space} (\text{bernoulli-stream } p)$ .  $w !! i$ } = p
and measure-bernoulli-stream-neg-snth: prob p { $w \in \text{space} (\text{bernoulli-stream } p)$ .  $\neg w !! i$ } =  $1 - p$ 
using measure-bernoulli-stream-snth-pred[OF assms, of {i}  $\lambda x. \text{True}$ ]
measure-bernoulli-stream-snth-pred[OF assms, of {i}  $\lambda x. \text{False}$ ] by auto

```

Now we can express the expected value of a single coin toss.

### lemma integral-toss:

```

assumes  $0 \leq p$   $p \leq 1$ 

```

```

shows expectation p (toss n) = 2 * p - 1
proof -
have [simp]:{s. s !! n} ∈ events p using measurable-snth[THEN measurable-sets,
of {True} measure-pmf (bernoulli-pmf p) n, folded bernoulli-stream-def]
by (simp add: vimage-def)
have expectation p (toss n) = Bochner-Integration.simple-bochner-integral (bernoulli-stream
p) (toss n)
using toss.random-variable[of n, THEN measurable-sets]
by (intro simple-bochner-integrable-eq-integral[symmetric] simple-bochner-integrable.intros)
(auto simp add: toss-def simple-function-def image-def)
also have ... = p - prob p {s. ¬ s !! n} unfolding simple-bochner-integral-def
using measure-bernoulli-stream-snth[OF assms]
by (simp add: range-toss, simp add: toss-def)
also have ... = p - (1 - prob p {s. s !! n}) by (subst prob-compl[symmetric],
auto simp add: Collect-neg-eq Compl-eq-Diff-UNIV)
finally show ?thesis using measure-bernoulli-stream-snth[OF assms] by simp
qed

```

Now, we show that the tosses are independent from one another.

```

lemma indep-vars-toss:
assumes 0 ≤ p p ≤ 1
shows indep-vars p (λ-. borel) toss {0..}
proof (subst indep-vars-def, intro conjI indep-sets-sigma)
{
fix A J assume asm: J ≠ {} finite J ∀j∈J. A j ∈ {toss j -` A ∩ space
(bernoulli-stream p) |A. A ∈ borel}
hence ∀j∈J. ∃B ∈ borel. A j = toss j -` B ∩ space (bernoulli-stream p) by
auto
then obtain B where B-is: A j = toss j -` B ∩ space (bernoulli-stream p)
B j ∈ borel if j ∈ J for j by metis

have prob p (⋂ (A ` J)) = (∏j∈J. prob p (A j))
proof cases

```

We consider the case where there is a zero probability event.

```

assume ∃j ∈ J. 1 ∉ B j ∧ -1 ∉ B j
then obtain j where j-is: j ∈ J 1 ∉ B j -1 ∉ B j by blast
hence A-j-empty: A j = {} using B-is by (force simp add: toss-def vimage-def)
hence ⋂ (A ` J) = {} using j-is by blast
moreover have prob p (A j) = 0 using A-j-empty by simp
ultimately show ?thesis using j-is asm(2) by auto
next

```

We now assume all events have positive probability.

```

assume ¬(∃j ∈ J. 1 ∉ B j ∧ -1 ∉ B j)
hence *: 1 ∈ B j ∨ -1 ∈ B j if j ∈ J for j using that by blast

```

```

define J' where [simp]: J' = {j ∈ J. (1 ∈ B j) ↔ (-1 ∉ B j)}

```

```

  hence toss j w ∈ B j ↔ (1 ∈ B j) = w !! j if j ∈ J' for w j using that
  unfolding toss-def by simp
    hence  $(\bigcap (A \setminus J')) = \{w \in \text{space}(\text{bernoulli-stream } p). \forall j \in J'. (1 \in B j) = w !! j\}$  using B-is by force
      hence prob-J': prob p  $(\bigcap (A \setminus J')) = p \wedge \text{card}(J' \cap \{j. 1 \in B j\}) * (1 - p) \wedge \text{card}(J' - \{j. 1 \in B j\})$ 
        using measure-bernoulli-stream-snth-pred[OF assms finite-subset[OF -asm(2)], of J' λj. 1 ∈ B j] by auto

The index set J' consists of the indices of all non-trivial events.

have A-j-True: A j =  $\{w \in \text{space}(\text{bernoulli-stream } p). w !! j\}$  if j ∈ J' ∩ {j. 1 ∈ B j} for j
  using that by (auto simp add: toss-def B-is(1) split: if-splits)

have A-j-False: A j =  $\{w \in \text{space}(\text{bernoulli-stream } p). \neg w !! j\}$  if j ∈ J' - {j. 1 ∈ B j} for j
  using that B-is by (auto simp add: toss-def)

have A-j-top: A j = space(bernoulli-stream p) if j ∈ J - J' for j using that
* by (auto simp add: B-is toss-def)
  hence  $\bigcap (A \setminus J) = \bigcap (A \setminus J')$  by auto
  hence prob p  $(\bigcap (A \setminus J)) = \text{prob } p (\bigcap (A \setminus J'))$  by presburger
  also have ... =  $(\prod_{j \in J'} \cap \{j. 1 \in B j\}. \text{prob } p (A j)) * (\prod_{j \in J' - \{j. 1 \in B j\}} \text{prob } p (A j))$ 
    by (simp only: prob-J' A-j-True A-j-False measure-bernoulli-stream-snth[OF assms] measure-bernoulli-stream-neg-snth[OF assms] cong: prod.cong) simp
    also have ... =  $(\prod_{j \in J'}. \text{prob } p (A j))$  using asm(2) by (intro prod.Int-Diff[symmetric])
  auto
  also have ... =  $(\prod_{j \in J'}. \text{prob } p (A j)) * (\prod_{j \in J - J'}. \text{prob } p (A j))$  using
A-j-top prob-space by simp
  also have ... =  $(\prod_{j \in J}. \text{prob } p (A j))$  using asm(2) by (metis (no-types,
  lifting) J'-def mem-Collect-eq mult.commute prod.subset-diff subsetI)
  finally show ?thesis .
qed
}

thus indep-sets p ( $\lambda i. \{\text{toss } i - ' A \cap \text{space}(\text{bernoulli-stream } p) | A. A \in \text{sets borel}\}$ ) {0..} using measurable-sets[OF toss.random-variable]
  by (intro indep-setsI subsetI) fastforce
qed (simp, intro Int-stableI, simp, metis sets.Int vimage-Int)

```

The fortune of a player is a martingale (resp. sub- or supermartingale) with respect to the filtration generated by the coin tosses.

```

theorem fortune-martingale:
  assumes p = 1/2
  shows martingale (bernoulli-stream p) (toss-filtration p) 0 fortune
  using cond-exp-indep[OF bernoulli-stream-natural-filtration.subalg indep-set-natural-filtration integrable-toss, OF zero-order(1) lessI indep-vars-toss, of p]
  integral-toss assms

```

**by** (*intro bernoulli-stream-natural-filtration.martingale-of-cond-exp-diff-Suc-eq-zero integrable-fortune fortune.adapted-process-axioms*) (*force simp add: fortune-toss-sum*)

**theorem** *fortune-submartingale*:

**assumes**  $1/2 \leq p$   $p \leq 1$

**shows** *submartingale* (*bernoulli-stream p*) (*toss-filtration p*) 0 *fortune*

**proof** (*intro bernoulli-stream-natural-filtration.submartingale-of-cond-exp-diff-Suc-nonneg integrable-fortune fortune.adapted-process-axioms*)

**fix** *n*

**show**  $\text{AE } \xi \text{ in bernoulli-stream } p. 0 \leq \text{cond-exp}(\text{bernoulli-stream } p) (\text{toss-filtration } p \text{ } n) (\lambda \xi. \text{fortune}(\text{Suc } n) \xi - \text{fortune } n \xi) \xi$

**using** *cond-exp-indep*[*OF bernoulli-stream-natural-filtration.subalg indep-set-natural-filtration integrable-toss, OF zero-order(1) lessI indep-vars-toss, of p n*]

*integral-toss*[*of p Suc n assms*

**by** (*force simp add: fortune-toss-sum*)

**qed**

**theorem** *fortune-supermartingale*:

**assumes**  $0 \leq p$   $p \leq 1/2$

**shows** *supermartingale* (*bernoulli-stream p*) (*toss-filtration p*) 0 *fortune*

**proof** (*intro bernoulli-stream-natural-filtration.supermartingale-of-cond-exp-diff-Suc-le-zero integrable-fortune fortune.adapted-process-axioms*)

**fix** *n*

**show**  $\text{AE } \xi \text{ in bernoulli-stream } p. 0 \geq \text{cond-exp}(\text{bernoulli-stream } p) (\text{toss-filtration } p \text{ } n) (\lambda \xi. \text{fortune}(\text{Suc } n) \xi - \text{fortune } n \xi) \xi$

**using** *cond-exp-indep*[*OF bernoulli-stream-natural-filtration.subalg indep-set-natural-filtration integrable-toss, OF zero-order(1) lessI indep-vars-toss, of p n*]

*integral-toss*[*of p Suc n assms*

**by** (*force simp add: fortune-toss-sum*)

**qed**

**end**

## References

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