

Lovasz Local Lemma

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Abstract

This entry aims to formalise several useful general techniques for using the *probabilistic method* for combinatorial structures (or discrete spaces more generally). In particular, it focuses on bounding tools, such as the union and complete independence bounds, and the first formalisation of the pivotal Lovász local lemma. The formalisation focuses on the general lemma, however also proves several useful variations, including the more well known symmetric version. Both the original formalisation and several of the variations used dependency graphs, which were formalised using Noschinski’s general directed graph library [2]. Additionally, the entry provides several useful existence lemmas, required at the end of most probabilistic proofs on combinatorial structures. Finally, the entry includes several significant extensions to the existing probability libraries, particularly for conditional probability (such as Bayes theorem) and independent events. The formalisation is primarily based on Alon and Spencer’s textbook [1], as well as Zhao’s course notes [3].

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1 Extensional function extras

Counting lemmas (i.e. reasoning on cardinality) of sets on the extensional function relation

```
theory PiE-Rel-Extras imports Card-Partitions.Card-Partitions
begin
```

1.1 Relations and Extensional Function sets

A number of lemmas to convert between relations and functions for counting purposes. Note, ultimately not needed in this formalisation, but may be of use in the future

```
lemma Range-unfold: Range  $r = \{y. \exists x. (x, y) \in r\}$ 
  <proof>
```

```
definition fun-to-rel:: 'a set  $\Rightarrow$  'b set  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a  $\times$  'b) set where
fun-to-rel A B f  $\equiv \{(a, b) \mid a \ b . a \in A \wedge b \in B \wedge f a = b\}$ 
```

```
definition rel-to-fun:: ('a  $\times$  'b) set  $\Rightarrow$  ('a  $\Rightarrow$  'b) where
rel-to-fun R  $\equiv \lambda a .$  (if a  $\in$  Domain R then (THE b . (a, b)  $\in$  R) else undefined)
```

```
lemma fun-to-relI:  $a \in A \Longrightarrow b \in B \Longrightarrow f a = b \Longrightarrow (a, b) \in$  fun-to-rel A B f
  <proof>
```

```
lemma fun-to-rel-alt: fun-to-rel A B f  $\equiv \{(a, f a) \mid a \ b . a \in A \wedge f a \in B\}$ 
  <proof>
```

```
lemma fun-to-relI2:  $a \in A \Longrightarrow f a \in B \Longrightarrow (a, f a) \in$  fun-to-rel A B f
  <proof>
```

lemma *rel-to-fun-in[simp]*: $a \in \text{Domain } R \implies (\text{rel-to-fun } R) a = (\text{THE } b . (a, b) \in R)$

<proof>

lemma *rel-to-fun-undefined[simp]*: $a \notin \text{Domain } R \implies (\text{rel-to-fun } R) a = \text{undefined}$

<proof>

lemma *single-valued-unique-Dom-iff*: $\text{single-valued } R \iff (\forall x \in \text{Domain } R. \exists! y . (x, y) \in R)$

<proof>

lemma *rel-to-fun-range*:

assumes *single-valued* R

assumes $a \in \text{Domain } R$

shows $(\text{THE } b . (a, b) \in R) \in \text{Range } R$

<proof>

lemma *rel-to-fun-extensional*: $\text{single-valued } R \implies \text{rel-to-fun } R \in (\text{Domain } R \rightarrow_E \text{Range } R)$

<proof>

lemma *single-value-fun-to-rel*: $\text{single-valued } (\text{fun-to-rel } A B f)$

<proof>

lemma *fun-to-rel-domain*:

assumes $f \in A \rightarrow_E B$

shows $\text{Domain } (\text{fun-to-rel } A B f) = A$

<proof>

lemma *fun-to-rel-range*:

assumes $f \in A \rightarrow_E B$

shows $\text{Range } (\text{fun-to-rel } A B f) \subseteq B$

<proof>

lemma *rel-to-fun-to-rel*:

assumes $f \in A \rightarrow_E B$

shows $\text{rel-to-fun } (\text{fun-to-rel } A B f) = f$

<proof>

lemma *fun-to-rel-to-fun*:

assumes *single-valued* R

shows $\text{fun-to-rel } (\text{Domain } R) (\text{Range } R) (\text{rel-to-fun } R) = R$

<proof>

lemma *bij-betw-fun-to-rel*:

assumes $f \in A \rightarrow_E B$

shows *bij-betw* $(\lambda a . (a, f a)) A (\text{fun-to-rel } A B f)$

<proof>

lemma *fun-to-rel-indiv-card*:
assumes $f \in A \rightarrow_E B$
shows $\text{card } (\text{fun-to-rel } A B f) = \text{card } A$
 $\langle \text{proof} \rangle$

lemma *fun-to-rel-inj*:
assumes $C \subseteq A \rightarrow_E B$
shows *inj-on* $(\text{fun-to-rel } A B) C$
 $\langle \text{proof} \rangle$

lemma *fun-to-rel-ss*: $\text{fun-to-rel } A B f \subseteq A \times B$
 $\langle \text{proof} \rangle$

lemma *card-fun-to-rel*: $C \subseteq A \rightarrow_E B \implies \text{card } C = \text{card } ((\lambda f . \text{fun-to-rel } A B f) ' C)$
 $\langle \text{proof} \rangle$

1.2 Cardinality Lemmas

Lemmas to count variations of filtered sets over the extensional function set relation

lemma *card-PiE-filter-range-set*:
assumes $\bigwedge a. a \in A' \implies X a \in C$
assumes $A' \subseteq A$
assumes *finite* A
shows $\text{card } \{f \in A \rightarrow_E C . \forall a \in A' . f a = X a\} = (\text{card } C) \frown (\text{card } A - \text{card } A')$
 $\langle \text{proof} \rangle$

lemma *card-PiE-filter-range-indiv*: $X a' \in C \implies a' \in A \implies \text{finite } A \implies$
 $\text{card } \{f \in A \rightarrow_E C . f a' = X a'\} = (\text{card } C) \frown (\text{card } A - 1)$
 $\langle \text{proof} \rangle$

lemma *card-PiE-filter-range-set-const*: $c \in C \implies A' \subseteq A \implies \text{finite } A \implies$
 $\text{card } \{f \in A \rightarrow_E C . \forall a \in A' . f a = c\} = (\text{card } C) \frown (\text{card } A - \text{card } A')$
 $\langle \text{proof} \rangle$

lemma *card-PiE-filter-range-set-nat*: $c \in \{0..<n\} \implies A' \subseteq A \implies \text{finite } A \implies$
 $\text{card } \{f \in A \rightarrow_E \{0..<n\} . \forall a \in A' . f a = c\} = n \frown (\text{card } A - \text{card } A')$
 $\langle \text{proof} \rangle$

end

2 Digraph extensions

Extensions to the existing library for directed graphs, basically neighborhood

theory *Digraph-Extensions*
imports

```

    Graph-Theory.Digraph
    Graph-Theory.Pair-Digraph
begin

definition (in pre-digraph) neighborhood :: 'a ⇒ 'a set where
neighborhood u ≡ {v ∈ verts G . dominates G u v}

lemma (in wf-digraph) neighborhood-wf: neighborhood v ⊆ verts G
⟨proof⟩

lemma (in pair-pre-digraph) neighborhood-alt:
neighborhood u = {v ∈ pverts G . (u, v) ∈ arcs G}
⟨proof⟩

lemma (in fin-digraph) neighborhood-finite: finite (neighborhood v)
⟨proof⟩

lemma (in wf-digraph) neighborhood-edge-iff: y ∈ neighborhood x ⟷ (x, y) ∈
arcs-ends G
⟨proof⟩

lemma (in loopfree-digraph) neighborhood-self-not: v ∉ (neighborhood v)
⟨proof⟩

lemma (in nomulti-digraph) inj-on-head-out-arcs: inj-on (head G) (out-arcs G u)
⟨proof⟩

lemma (in nomulti-digraph) out-degree-neighborhood: out-degree G u = card (neighborhood
u)
⟨proof⟩

lemma (in digraph) neighborhood-empty-iff: out-degree G u = 0 ⟷ neighborhood
u = {}
⟨proof⟩

end

```

3 General Event Lemmas

General lemmas for reasoning on events in probability spaces after different operations

```

theory Prob-Events-Extras
imports
    HOL-Probability.Probability
    PiE-Rel-Extras
begin

context prob-space

```

begin

lemma *prob-sum-Union:*

assumes *measurable: finite A A ⊆ events disjoint A*

shows $\text{prob} (\bigcup A) = (\sum_{e \in A. \text{prob} (e)})$

<proof>

lemma *events-inter:*

assumes *finite S*

assumes $S \neq \{\}$

shows $(\bigwedge A. A \in S \implies A \in \text{events}) \implies \bigcap S \in \text{events}$

<proof>

lemma *events-union:*

assumes *finite S*

shows $(\bigwedge A. A \in S \implies A \in \text{events}) \implies \bigcup S \in \text{events}$

<proof>

lemma *prob-inter-set-lt-lem:* $A \in \text{events} \implies \text{prob} (A \cap (\bigcap AS)) \leq \text{prob} A$

<proof>

lemma *Inter-event-ss:* $\text{finite } A \implies A \subseteq \text{events} \implies A \neq \{\} \implies \bigcap A \in \text{events}$

<proof>

lemma *prob-inter-ss-lt:*

assumes *finite A*

assumes $A \subseteq \text{events}$

assumes $B \neq \{\}$

assumes $B \subseteq A$

shows $\text{prob} (\bigcap A) \leq \text{prob} (\bigcap B)$

<proof>

lemma *prob-inter-ss-lt-index:*

assumes *finite A*

assumes $F \text{ ' } A \subseteq \text{events}$

assumes $B \neq \{\}$

assumes $B \subseteq A$

shows $\text{prob} (\bigcap (F \text{ ' } A)) \leq \text{prob} (\bigcap (F \text{ ' } B))$

<proof>

lemma *space-compl-double:*

assumes $S \subseteq \text{events}$

shows $((-) (\text{space } M)) \text{ ' } (((-) (\text{space } M)) \text{ ' } S) = S$

<proof>

lemma *bij-betw-compl-sets:*

assumes $S \subseteq \text{events}$

assumes $S' = ((-) (\text{space } M)) \text{ ' } S$

shows *bij-betw* $((-) (\text{space } M)) S' S$

<proof>

lemma *bij-betw-compl-sets-rev*:

assumes $S \subseteq \text{events}$

assumes $S' = ((-) (\text{space } M)) ' S$

shows *bij-betw* $((-) (\text{space } M)) S S'$

<proof>

lemma *prob0-basic-inter*: $A \in \text{events} \implies B \in \text{events} \implies \text{prob } A = 0 \implies \text{prob } (A \cap B) = 0$

<proof>

lemma *prob0-basic-Inter*: $A \in \text{events} \implies B \subseteq \text{events} \implies \text{prob } A = 0 \implies \text{prob } (A \cap (\bigcap B)) = 0$

<proof>

lemma *prob1-basic-inter*: $A \in \text{events} \implies B \in \text{events} \implies \text{prob } A = 1 \implies \text{prob } (A \cap B) = \text{prob } B$

<proof>

lemma *prob1-basic-Inter*:

assumes $A \in \text{events}$ $B \subseteq \text{events}$

assumes $\text{prob } A = 1$

assumes $B \neq \{\}$

assumes *finite* B

shows $\text{prob } (A \cap (\bigcap B)) = \text{prob } (\bigcap B)$

<proof>

lemma *compl-identity*: $A \in \text{events} \implies \text{space } M - (\text{space } M - A) = A$

<proof>

lemma *prob-addition-rule*: $A \in \text{events} \implies B \in \text{events} \implies$

$\text{prob } (A \cup B) = \text{prob } A + \text{prob } B - \text{prob } (A \cap B)$

<proof>

lemma *compl-subset-in-events*: $S \subseteq \text{events} \implies (-) (\text{space } M) ' S \subseteq \text{events}$

<proof>

lemma *prob-compl-diff-inter*: $A \in \text{events} \implies B \in \text{events} \implies$

$\text{prob } (A \cap (\text{space } M - B)) = \text{prob } A - \text{prob } (A \cap B)$

<proof>

lemma *bij-betw-prod-prob*: *bij-betw* $f A B \implies (\prod_{b \in B}. \text{prob } b) = (\prod_{a \in A}. \text{prob } (f a))$

<proof>

definition *event-compl* :: 'a set \implies 'a set **where**

event-compl $A \equiv \text{space } M - A$

lemma compl-Union: $A \neq \{\}$ \implies $\text{space } M - (\bigcup A) = (\bigcap a \in A . (\text{space } M - a))$

<proof>

lemma compl-Union-fn: $A \neq \{\}$ \implies $\text{space } M - (\bigcup (F \cdot A)) = (\bigcap a \in A . (\text{space } M - F a))$

<proof>

end

Reasoning on the probability of function sets

lemma card-PiE-val-ss-eq:

assumes *finite* A

assumes $b \in B$

assumes $d \subseteq A$

assumes $B \neq \{\}$

assumes *finite* B

shows $\text{card } \{f \in (A \rightarrow_E B) . (\forall v \in d . f v = b)\} / \text{card } (A \rightarrow_E B) = 1 / ((\text{card } B) \text{ powi } (\text{card } d))$

(is $\text{card } \{f \in ?C . (\forall v \in d . f v = b)\} / \text{card } ?C = 1 / ((\text{card } B) \text{ powi } (\text{card } d))$ **)**

<proof>

lemma card-PiE-val-indiv-eq:

assumes *finite* A

assumes $b \in B$

assumes $d \in A$

assumes $B \neq \{\}$

assumes *finite* B

shows $\text{card } \{f \in (A \rightarrow_E B) . f d = b\} / \text{card } (A \rightarrow_E B) = 1 / (\text{card } B)$

(is $\text{card } \{f \in ?C . f d = b\} / \text{card } ?C = 1 / (\text{card } B)$ **)**

<proof>

lemma prob-uniform-ex-fun-space:

assumes *finite* A

assumes $b \in B$

assumes $d \subseteq A$

assumes $B \neq \{\}$

assumes $A \neq \{\}$

assumes *finite* B

shows $\text{prob-space.prob } (\text{uniform-count-measure } (A \rightarrow_E B)) \{f \in (A \rightarrow_E B) . (\forall v \in d . f v = b)\} =$

$1 / ((\text{card } B) \text{ powi } (\text{card } d))$

<proof>

proposition integrable-uniform-count-measure-finite:

fixes $g :: 'a \Rightarrow 'b :: \{\text{banach, second-countable-topology}\}$

shows *finite* $A \implies$ *integrable* $(\text{uniform-count-measure } A) g$

<proof>

end

4 Conditional Probability Library Extensions

```
theory Cond-Prob-Extensions
  imports
    Prob-Events-Extras
    Design-Theory.Multisets-Extras
begin
```

4.1 Miscellaneous Set and List Lemmas

```
lemma nth-image-tl:
  assumes  $xs \neq []$ 
  shows  $nth\ xs\ ' \{1..<length\ xs\} = set(tl\ xs)$ 
<proof>
```

```
lemma exists-list-card:
  assumes finite S
  obtains xs where  $set\ xs = S$  and  $length\ xs = card\ S$ 
<proof>
```

```
lemma bij-betw-inter-empty:
  assumes bij-betw f A B
  assumes  $A' \subseteq A$ 
  assumes  $A'' \subseteq A$ 
  assumes  $A' \cap A'' = \{\}$ 
  shows  $f\ ' \ A' \cap f\ ' \ A'' = \{\}$ 
<proof>
```

```
lemma bij-betw-image-comp-eq:
  assumes bij-betw g T S
  shows  $(F \circ g)\ ' \ T = F\ ' \ S$ 
<proof>
```

```
lemma prod-card-image-set-eq:
  assumes bij-betw f {0..<card S} S
  assumes finite S
  shows  $(\prod i \in \{n..<(card\ S)\} . g\ (f\ i)) = (\prod i \in f\ ' \ \{n..<card\ S\} . g\ i)$ 
<proof>
```

```
lemma set-take-distinct-elem-not:
  assumes distinct xs
  assumes  $i < length\ xs$ 
  shows  $xs\ !\ i \notin set\ (take\ i\ xs)$ 
<proof>
```

4.2 Conditional Probability Basics

context *prob-space*

begin

Abbreviation to mirror mathematical notations

abbreviation *cond-prob-ev* :: 'a set \Rightarrow 'a set \Rightarrow real ($\langle \mathcal{P}'(- | -) \rangle$) **where**
 $\mathcal{P}(B | A) \equiv \mathcal{P}(x \text{ in } M. (x \in B) | (x \in A))$

lemma *cond-prob-inter*: $\mathcal{P}(B | A) = \mathcal{P}(\omega \text{ in } M. (\omega \in B \cap A)) / \mathcal{P}(\omega \text{ in } M. (\omega \in A))$
\langle proof \rangle

lemma *cond-prob-ev-def*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(B | A) = \text{prob } (A \cap B) / \text{prob } A$
\langle proof \rangle

lemma *measurable-in-ev*:
assumes $A \in \text{events}$
shows $\text{Measurable.pred } M (\lambda x . x \in A)$
\langle proof \rangle

lemma *measure-uniform-measure-eq-cond-prob-ev*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A | B) = \mathcal{P}(x \text{ in uniform-measure } M \{x \in \text{space } M. x \in B\}. x \in A)$
\langle proof \rangle

lemma *measure-uniform-measure-eq-cond-prob-ev2*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A | B) = \text{measure } (\text{uniform-measure } M \{x \in \text{space } M. x \in B\}) A$
\langle proof \rangle

lemma *measure-uniform-measure-eq-cond-prob-ev3*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A | B) = \text{measure } (\text{uniform-measure } M B) A$
\langle proof \rangle

lemma *prob-space-cond-prob-uniform*:
assumes $\text{prob } (\{x \in \text{space } M. Q x\}) > 0$
shows $\text{prob-space } (\text{uniform-measure } M \{x \in \text{space } M. Q x\})$
\langle proof \rangle

lemma *prob-space-cond-prob-event*:
assumes $\text{prob } B > 0$
shows $\text{prob-space } (\text{uniform-measure } M B)$
\langle proof \rangle

Note this case shouldn't be used. Conditional probability should have > 0 assumption

lemma *cond-prob-empty*: $\mathcal{P}(B \mid \{\}) = 0$
<proof>

lemma *cond-prob-space*: $\mathcal{P}(A \mid \text{space } M) = \mathcal{P}(w \text{ in } M . w \in A)$
<proof>

lemma *cond-prob-space-ev*: **assumes** $A \in \text{events}$ **shows** $\mathcal{P}(A \mid \text{space } M) = \text{prob } A$
<proof>

lemma *cond-prob-UNIV*: $\mathcal{P}(A \mid \text{UNIV}) = \mathcal{P}(w \text{ in } M . w \in A)$
<proof>

lemma *cond-prob-UNIV-ev*: $A \in \text{events} \implies \mathcal{P}(A \mid \text{UNIV}) = \text{prob } A$
<proof>

lemma *cond-prob-neg*:
assumes $A \in \text{events } B \in \text{events}$
assumes $\text{prob } A > 0$
shows $\mathcal{P}(\text{space } M - B \mid A) = 1 - \mathcal{P}(B \mid A)$
<proof>

4.3 Bayes Theorem

lemma *prob-intersect-A*:
assumes $A \in \text{events } B \in \text{events}$
shows $\text{prob } (A \cap B) = \text{prob } A * \mathcal{P}(B \mid A)$
<proof>

lemma *prob-intersect-B*:
assumes $A \in \text{events } B \in \text{events}$
shows $\text{prob } (A \cap B) = \text{prob } B * \mathcal{P}(A \mid B)$
<proof>

theorem *Bayes-theorem*:
assumes $A \in \text{events } B \in \text{events}$
shows $\text{prob } B * \mathcal{P}(A \mid B) = \text{prob } A * \mathcal{P}(B \mid A)$
<proof>

corollary *Bayes-theorem-div*:
assumes $A \in \text{events } B \in \text{events}$
shows $\mathcal{P}(A \mid B) = (\text{prob } A * \mathcal{P}(B \mid A)) / (\text{prob } B)$
<proof>

lemma *cond-prob-dual-intersect*:
assumes $A \in \text{events } B \in \text{events } C \in \text{events}$
assumes $\text{prob } C \neq 0$
shows $\mathcal{P}(A \mid (B \cap C)) = \mathcal{P}(A \cap B \mid C) / \mathcal{P}(B \mid C)$ (is ?LHS = ?RHS)
<proof>

lemma *cond-prob-ev-double*:

assumes $A \in \text{events}$ $B \in \text{events}$ $C \in \text{events}$

assumes $\text{prob } C > 0$

shows $\mathcal{P}(x \text{ in } (\text{uniform-measure } M \ C). (x \in A) \mid (x \in B)) = \mathcal{P}(A \mid (B \cap C))$

<proof>

lemma *cond-prob-inter-set-lt*:

assumes $A \in \text{events}$ $B \in \text{events}$ $AS \subseteq \text{events}$

assumes *finite* AS

shows $\mathcal{P}((A \cap (\bigcap AS)) \mid B) \leq \mathcal{P}(A \mid B)$ (**is** ?LHS \leq ?RHS)

<proof>

4.4 Conditional Probability Multiplication Rule

Many list and indexed variations of this lemma

lemma *prob-cond-Inter-List*:

assumes $xs \neq []$

assumes $\bigwedge A. A \in \text{set } xs \implies A \in \text{events}$

shows $\text{prob } (\bigcap (\text{set } xs)) = \text{prob } (\text{hd } xs) * (\prod i = 1..<(\text{length } xs) .$

$\mathcal{P}((xs ! i) \mid (\bigcap (\text{set } (\text{take } i \ xs))))$

<proof>

lemma *prob-cond-Inter-index*:

fixes $n :: \text{nat}$

assumes $n > 0$

assumes $F \text{ ' } \{0..<n\} \subseteq \text{events}$

shows $\text{prob } (\bigcap (F \text{ ' } \{0..<n\})) = \text{prob } (F \ 0) * (\prod i \in \{1..<n\} .$

$\mathcal{P}(F \ i \mid (\bigcap (F \text{ ' } \{0..<i\})))$

<proof>

lemma *prob-cond-Inter-index-compl*:

fixes $n :: \text{nat}$

assumes $n > 0$

assumes $F \text{ ' } \{0..<n\} \subseteq \text{events}$

shows $\text{prob } (\bigcap x \in \{0..<n\} . \text{space } M - F \ x) = \text{prob } (\text{space } M - F \ 0) * (\prod i$

$\in \{1..<n\} .$

$\mathcal{P}(\text{space } M - F \ i \mid (\bigcap j \in \{0..<i\} . \text{space } M - F \ j)))$

<proof>

lemma *prob-cond-Inter-take-cond*:

assumes $xs \neq []$

assumes $\text{set } xs \subseteq \text{events}$

assumes $S \subseteq \text{events}$

assumes $S \neq \{\}$

assumes *finite* S

assumes $\text{prob } (\bigcap S) > 0$

shows $\mathcal{P}((\bigcap(\text{set } xs)) \mid (\bigcap S)) = (\prod i = 0..<(\text{length } xs) . \mathcal{P}((xs ! i) \mid (\bigcap(\text{set } (\text{take } i \text{ } xs) \cup S))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-index-cond-set:*

fixes $n :: \text{nat}$
assumes $n > 0$
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F \text{ ' } \{0..<n\} \subseteq \text{events}$
assumes $\text{prob } (\bigcap E) > 0$
shows $\mathcal{P}((\bigcap(F \text{ ' } \{0..<n\})) \mid (\bigcap E)) = (\prod i \in \{0..<n\} . \mathcal{P}(F i \mid (\bigcap((F \text{ ' } \{0..<i\}) \cup E))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-index-cond-compl-set:*

fixes $n :: \text{nat}$
assumes $n > 0$
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F \text{ ' } \{0..<n\} \subseteq \text{events}$
assumes $\text{prob } (\bigcap E) > 0$
shows $\mathcal{P}((\bigcap((-) (\text{space } M) \text{ ' } F \text{ ' } \{0..<n\})) \mid (\bigcap E)) = (\prod i = 0..<n . \mathcal{P}((\text{space } M - F i) \mid (\bigcap((-) (\text{space } M) \text{ ' } F \text{ ' } \{0..<i\} \cup E))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-index-cond:*

fixes $n :: \text{nat}$
assumes $n > 0$
assumes $n < m$
assumes $F \text{ ' } \{0..<m\} \subseteq \text{events}$
assumes $\text{prob } (\bigcap j \in \{n..<m\} . F j) > 0$
shows $\mathcal{P}((\bigcap(F \text{ ' } \{0..<n\})) \mid (\bigcap j \in \{n..<m\} . F j)) = (\prod i \in \{0..<n\} . \mathcal{P}(F i \mid (\bigcap((F \text{ ' } \{0..<i\}) \cup (F \text{ ' } \{n..<m\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-index-cond-compl:*

fixes $n :: \text{nat}$
assumes $n > 0$
assumes $n < m$
assumes $F \text{ ' } \{0..<m\} \subseteq \text{events}$
assumes $\text{prob } (\bigcap j \in \{n..<m\} . F j) > 0$
shows $\mathcal{P}((\bigcap((-) (\text{space } M) \text{ ' } F \text{ ' } \{0..<n\})) \mid (\bigcap(F \text{ ' } \{n..<m\}))) = (\prod i = 0..<n . \mathcal{P}((\text{space } M - F i) \mid (\bigcap((-) (\text{space } M) \text{ ' } F \text{ ' } \{0..<i\} \cup (F \text{ ' } \{n..<m\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-take-cond-neg*:

assumes $xs \neq []$
assumes $set\ xs \subseteq events$
assumes $S \subseteq events$
assumes $S \neq \{\}$
assumes *finite* S
assumes $prob\ (\bigcap S) > 0$
shows $\mathcal{P}((\bigcap((-)\ (space\ M)\ '(\set\ xs))) \mid (\bigcap S)) =$
 $(\prod_{i=0..<(length\ xs)} . \mathcal{P}((space\ M - xs\ !\ i) \mid (\bigcap((-)\ (space\ M)\ '(\set\ (take\ i\ xs)) \cup S))))$
<proof>

lemma *prob-cond-Inter-List-Index*:

assumes $xs \neq []$
assumes $set\ xs \subseteq events$
shows $prob\ (\bigcap(\set\ xs)) = prob\ (hd\ xs) * (\prod_{i=1..<(length\ xs)} .$
 $\mathcal{P}((xs\ !\ i) \mid (\bigcap_{j \in \{0..<i\}} . xs\ !\ j)))$
<proof>

lemma *obtains-prob-cond-Inter-index*:

assumes $S \neq \{\}$
assumes $S \subseteq events$
assumes *finite* S
obtains xs **where** $set\ xs = S$ **and** $length\ xs = card\ S$ **and**
 $prob\ (\bigcap S) = prob\ (hd\ xs) * (\prod_{i=1..<(length\ xs)} . \mathcal{P}((xs\ !\ i) \mid (\bigcap_{j \in \{0..<i\}} . xs\ !\ j)))$
<proof>

lemma *obtain-list-index*:

assumes *bij-betw* $g\ \{0..<card\ S\}\ S$
assumes *finite* S
obtains xs **where** $set\ xs = S$ **and** $\bigwedge i . i \in \{0..<card\ S\} \implies g\ i = xs\ !\ i$ **and**
distinct xs
<proof>

lemma *prob-cond-inter-fn*:

assumes *bij-betw* $g\ \{0..<card\ S\}\ S$
assumes *finite* S
assumes $S \neq \{\}$
assumes $S \subseteq events$
shows $prob\ (\bigcap S) = prob\ (g\ 0) * (\prod_{i \in \{1..<(card\ S)\}} . \mathcal{P}(g\ i \mid (\bigcap_{j \in \{0..<i\}} (g\ j))))$
<proof>

lemma *prob-cond-inter-obtain-fn*:

assumes $S \neq \{\}$
assumes $S \subseteq events$
assumes *finite* S

obtains f where $\text{bij-betw } f \{0..<\text{card } S\} S$ and
 $\text{prob } (\bigcap S) = \text{prob } (f 0) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(f i \mid (\bigcap (f ' \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma $\text{prob-cond-inter-obtain-fn-compl}$:

assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
assumes $\text{finite } S$
obtains f where $\text{bij-betw } f \{0..<\text{card } S\} S$ and $\text{prob } (\bigcap ((-) (\text{space } M) ' S))$
 $=$
 $\text{prob } (\text{space } M - f 0) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(\text{space } M - f i \mid (\bigcap ((-) (\text{space } M) ' f ' \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma $\text{prob-cond-Inter-index-cond-fn}$:

assumes $I \neq \{\}$
assumes $\text{finite } I$
assumes $\text{finite } E$
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob } (\bigcap E) > 0$
assumes $\text{bb: } \text{bij-betw } g \{0..<\text{card } I\} I$
shows $\mathcal{P}((\bigcap (F ' g ' \{0..<\text{card } I\})) \mid (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\} . \mathcal{P}(F (g i) \mid (\bigcap ((F ' g ' \{0..<i\}) \cup E))))$
 $\langle \text{proof} \rangle$

lemma $\text{prob-cond-Inter-index-cond-obtains}$:

assumes $I \neq \{\}$
assumes $\text{finite } I$
assumes $\text{finite } E$
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob } (\bigcap E) > 0$
obtains g where $\text{bij-betw } g \{0..<\text{card } I\} I$ and $\mathcal{P}((\bigcap (F ' g ' \{0..<\text{card } I\})) \mid$
 $(\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\} . \mathcal{P}(F (g i) \mid (\bigcap ((F ' g ' \{0..<i\}) \cup E))))$
 $\langle \text{proof} \rangle$

lemma $\text{prob-cond-Inter-index-cond-compl-fn}$:

assumes $I \neq \{\}$
assumes $\text{finite } I$
assumes $\text{finite } E$
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob } (\bigcap E) > 0$

assumes $bb: \text{bij-betw } g \{0..<\text{card } I\} I$
shows $\mathcal{P}((\bigcap Aj \in I . \text{space } M - F Aj) \mid (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\}. \mathcal{P}(\text{space } M - F (g i) \mid (\bigcap ((\lambda Aj. \text{space } M - F Aj) ' g ' \{0..<i\}) \cup E))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-index-cond-compl-obtains:*

assumes $I \neq \{\}$
assumes *finite* I
assumes *finite* E
assumes $E \neq \{\}$
assumes $E \subseteq \text{events}$
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob } (\bigcap E) > 0$
obtains g **where** $\text{bij-betw } g \{0..<\text{card } I\} I$ **and** $\mathcal{P}((\bigcap Aj \in I . \text{space } M - F Aj) \mid (\bigcap E)) =$
 $(\prod i \in \{0..<\text{card } I\}. \mathcal{P}(\text{space } M - F (g i) \mid (\bigcap ((\lambda Aj. \text{space } M - F Aj) ' g ' \{0..<i\}) \cup E))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-index-fn2:*

assumes $F ' S \subseteq \text{events}$
assumes *finite* S
assumes $\text{card } S > 0$
assumes $\text{bij-betw } g \{0..<\text{card } S\} S$
shows $\text{prob } (\bigcap (F ' S)) = \text{prob } (F (g 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(F (g i) \mid (\bigcap (F ' g ' \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-index-fn:*

assumes $F ' S \subseteq \text{events}$
assumes *finite* S
assumes $S \neq \{\}$
assumes $\text{bij-betw } g \{0..<\text{card } S\} S$
shows $\text{prob } (\bigcap (F ' S)) = \text{prob } (F (g 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(F (g i) \mid (\bigcap (F ' g ' \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-index-obtain-fn:*

assumes $F ' S \subseteq \text{events}$
assumes *finite* S
assumes $S \neq \{\}$
obtains g **where** $\text{bij-betw } g \{0..<\text{card } S\} S$ **and**
 $\text{prob } (\bigcap (F ' S)) = \text{prob } (F (g 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(F (g i) \mid (\bigcap (F ' g ' \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-index-fn-compl:*

assumes $S \neq \{\}$

assumes $F \text{ ' } S \subseteq \text{events}$
assumes $\text{finite } S$
assumes $\text{bij-betw } f \{0..<\text{card } S\} S$
shows $\text{prob } (\bigcap ((-) (\text{space } M) \text{ ' } F \text{ ' } S)) = \text{prob } (\text{space } M - F (f 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(\text{space } M - F (f i) \mid (\bigcap ((-) (\text{space } M) \text{ ' } F \text{ ' } f \text{ ' } \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-inter-index-obtain-fn-compl:*

assumes $S \neq \{\}$
assumes $F \text{ ' } S \subseteq \text{events}$
assumes $\text{finite } S$
obtains f **where** $\text{bij-betw } f \{0..<\text{card } S\} S$ **and**
 $\text{prob } (\bigcap ((-) (\text{space } M) \text{ ' } F \text{ ' } S)) = \text{prob } (\text{space } M - F (f 0)) * (\prod i \in \{1..<(\text{card } S)\} . \mathcal{P}(\text{space } M - F (f i) \mid (\bigcap ((-) (\text{space } M) \text{ ' } F \text{ ' } f \text{ ' } \{0..<i\}))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-take:*

assumes $S \neq \{\}$
assumes $S \subseteq \text{events}$
assumes $\text{finite } S$
obtains xs **where** $\text{set } xs = S$ **and** $\text{length } xs = \text{card } S$ **and**
 $\text{prob } (\bigcap S) = \text{prob } (\text{hd } xs) * (\prod i = 1..<(\text{length } xs) . \mathcal{P}((xs ! i) \mid (\bigcap (\text{set } (\text{take } i xs))))))$
 $\langle \text{proof} \rangle$

lemma *prob-cond-Inter-set-bound:*

assumes $A \neq \{\}$
assumes $A \subseteq \text{events}$
assumes $\text{finite } A$
assumes $\bigwedge Ai . f Ai \geq 0 \wedge f Ai \leq 1$
assumes $\bigwedge Ai S . Ai \in A \implies S \subseteq A - \{Ai\} \implies S \neq \{\} \implies \mathcal{P}(Ai \mid (\bigcap S)) \geq f Ai$
assumes $\bigwedge Ai . Ai \in A \implies \text{prob } Ai \geq f Ai$
shows $\text{prob } (\bigcap A) \geq (\prod a' \in A . f a')$
 $\langle \text{proof} \rangle$
end

end

5 Independent Events

theory *Indep-Events* **imports** *Cond-Prob-Extensions*
begin

5.1 More bijection helpers

lemma *bij-betw-obtain-subset1*:

assumes *bij-betw* f A B

assumes $A' \subseteq A$

obtains B' **where** $B' \subseteq B$ **and** $B' = f \text{ ` } A'$

<proof>

lemma *bij-betw-obtain-subset2*:

assumes *bij-betw* f A B

assumes $B' \subseteq B$

obtains A' **where** $A' \subseteq A$ **and** $B' = f \text{ ` } A'$

<proof>

lemma *bij-betw-remove*: *bij-betw* f A $B \implies a \in A \implies \textit{bij-betw} f $(A - \{a\})$ $(B - \{f\ a\})$$

<proof>

5.2 Independent Event Extensions

Extensions on both the *indep_event* definition and the *indep_events* definition

context *prob-space*

begin

lemma *indep-eventsD*: *indep-events* A $I \implies (A \text{ ` } I \subseteq \textit{events}) \implies J \subseteq I \implies J \neq \{\}$ $\implies \textit{finite}$ $J \implies$

$\textit{prob} (\bigcap_{j \in J}. A\ j) = (\prod_{j \in J}. \textit{prob} (A\ j))$

<proof>

lemma

assumes *indep*: *indep-event* A B

shows *indep-eventD-ev1*: $A \in \textit{events}$

and *indep-eventD-ev2*: $B \in \textit{events}$

<proof>

lemma *indep-eventD*:

assumes *ie*: *indep-event* A B

shows $\textit{prob} (A \cap B) = \textit{prob} (A) * \textit{prob} (B)$

<proof>

lemma *indep-eventI[intro]*:

assumes *ev*: $A \in \textit{events}$ $B \in \textit{events}$

and *indep*: $\textit{prob} (A \cap B) = \textit{prob} A * \textit{prob} B$

shows *indep-event* A B

<proof>

Alternate set definition - when no possibility of duplicate objects

definition *indep-events-set* :: 'a set set \Rightarrow bool **where**

indep-events-set $E \equiv (E \subseteq \text{events} \wedge (\forall J. J \subseteq E \longrightarrow \text{finite } J \longrightarrow J \neq \{\} \longrightarrow \text{prob} (\bigcap J) = (\prod_{i \in J} \text{prob } i)))$

lemma *indep-events-setI[intro]*: $E \subseteq \text{events} \Longrightarrow (\bigwedge J. J \subseteq E \Longrightarrow \text{finite } J \Longrightarrow J \neq \{\} \Longrightarrow \text{prob} (\bigcap J) = (\prod_{i \in J} \text{prob } i)) \Longrightarrow \text{indep-events-set } E$
 <proof>

lemma *indep-events-subset*:
indep-events-set $E \longleftrightarrow (\forall J \subseteq E. \text{indep-events-set } J)$
 <proof>

lemma *indep-events-subset2*:
indep-events-set $E \Longrightarrow J \subseteq E \Longrightarrow \text{indep-events-set } J$
 <proof>

lemma *indep-events-set-events*: *indep-events-set* $E \Longrightarrow (\bigwedge e. e \in E \Longrightarrow e \in \text{events})$
 <proof>

lemma *indep-events-set-events-ss*: *indep-events-set* $E \Longrightarrow E \subseteq \text{events}$
 <proof>

lemma *indep-events-set-probs*: *indep-events-set* $E \Longrightarrow J \subseteq E \Longrightarrow \text{finite } J \Longrightarrow J \neq \{\} \Longrightarrow \text{prob} (\bigcap J) = (\prod_{i \in J} \text{prob } i)$
 <proof>

lemma *indep-events-set-prod-all*: *indep-events-set* $E \Longrightarrow \text{finite } E \Longrightarrow E \neq \{\} \Longrightarrow \text{prob} (\bigcap E) = \text{prod prob } E$
 <proof>

lemma *indep-events-not-contain-compl*:
assumes *indep-events-set* E
assumes $A \in E$
assumes $\text{prob } A > 0 \text{ prob } A < 1$
shows $(\text{space } M - A) \notin E$ (**is** $?A' \notin E$)
 <proof>

lemma *indep-events-contain-compl-prob01*:
assumes *indep-events-set* E
assumes $A \in E$
assumes $\text{space } M - A \in E$
shows $\text{prob } A = 0 \vee \text{prob } A = 1$
 <proof>

lemma *indep-events-set-singleton*:
assumes $A \in \text{events}$
shows *indep-events-set* $\{A\}$

<proof>

lemma *indep-events-pairs:*
 assumes *indep-events-set S*
 assumes $A \in S \ B \in S \ A \neq B$
 shows *indep-event A B*
 <proof>

lemma *indep-events-inter-pairs:*
 assumes *indep-events-set S*
 assumes *finite A finite B*
 assumes $A \neq \{\} \ B \neq \{\}$
 assumes $A \subseteq S \ B \subseteq S \ A \cap B = \{\}$
 shows *indep-event ($\bigcap A$) ($\bigcap B$)*
 <proof>

lemma *indep-events-inter-single:*
 assumes *indep-events-set S*
 assumes *finite B*
 assumes $B \neq \{\}$
 assumes $A \in S \ B \subseteq S \ A \notin B$
 shows *indep-event A ($\bigcap B$)*
 <proof>

lemma *indep-events-set-prob1:*
 assumes $A \in \text{events}$
 assumes $\text{prob } A = 1$
 assumes $A \notin S$
 assumes *indep-events-set S*
 shows *indep-events-set ($S \cup \{A\}$)*
 <proof>

lemma *indep-events-set-prob0:*
 assumes $A \in \text{events}$
 assumes $\text{prob } A = 0$
 assumes $A \notin S$
 assumes *indep-events-set S*
 shows *indep-events-set ($S \cup \{A\}$)*
 <proof>

lemma *indep-event-commute:*
 assumes *indep-event A B*
 shows *indep-event B A*
 <proof>

Showing complement operation maintains independence

lemma *indep-event-one-compl:*

assumes *indep-event* $A B$
shows *indep-event* A (*space* $M - B$)
 \langle *proof* \rangle

lemma *indep-event-one-compl-rev*:
assumes $B \in \text{events}$
assumes *indep-event* A (*space* $M - B$)
shows *indep-event* $A B$
 \langle *proof* \rangle

lemma *indep-event-double-compl*: *indep-event* $A B \implies \text{indep-event}$ (*space* $M - A$) (*space* $M - B$)
 \langle *proof* \rangle

lemma *indep-event-double-compl-rev*: $A \in \text{events} \implies B \in \text{events} \implies$
indep-event (*space* $M - A$) (*space* $M - B$) $\implies \text{indep-event}$ $A B$
 \langle *proof* \rangle

lemma *indep-events-set-one-compl*:
assumes *indep-events-set* S
assumes $A \in S$
shows *indep-events-set* ($\{\text{space } M - A\} \cup (S - \{A\})$)
 \langle *proof* \rangle

lemma *indep-events-set-update-compl*:
assumes *indep-events-set* E
assumes $E = A \cup B$
assumes $A \cap B = \{\}$
assumes *finite* E
shows *indep-events-set* ($((-) \text{space } M - A) \cup B$)
 \langle *proof* \rangle

lemma *indep-events-set-compl*:
assumes *indep-events-set* E
assumes *finite* E
shows *indep-events-set* ($(\lambda e. \text{space } M - e) \text{ ` } E$)
 \langle *proof* \rangle

lemma *indep-event-empty*:
assumes $A \in \text{events}$
shows *indep-event* $A \{\}$
 \langle *proof* \rangle

lemma *indep-event-compl-inter*:
assumes *indep-event* $A C$
assumes $B \in \text{events}$
assumes *indep-event* A ($B \cap C$)
shows *indep-event* A ($(\text{space } M - B) \cap C$)

<proof>

lemma *indep-events-index-subset:*

indep-events F E \longleftrightarrow ($\forall J \subseteq E. \text{indep-events F J}$)

<proof>

lemma *indep-events-index-subset2:*

indep-events F E \implies J \subseteq E \implies indep-events F J

<proof>

lemma *indep-events-events-ss: indep-events F E \implies F ' E \subseteq events*

<proof>

lemma *indep-events-events: indep-events F E \implies ($\bigwedge e. e \in E \implies F e \in \text{events}$)*

<proof>

lemma *indep-events-probs: indep-events F E \implies J \subseteq E \implies finite J \implies J \neq {}*

$\implies \text{prob} (\bigcap (F ' J)) = (\prod_{i \in J. \text{prob} (F i)}$)

<proof>

lemma *indep-events-prod-all: indep-events F E \implies finite E \implies E \neq {} \implies prob*

($\bigcap (F ' E)$) = ($\prod_{i \in E. \text{prob} (F i)$)

<proof>

lemma *indep-events-ev-not-contain-compl:*

assumes *indep-events F E*

assumes *A \in E*

assumes *prob (F A) > 0 prob (F A) < 1*

shows *(space M - F A) \notin F ' E (is ?A' \notin F ' E)*

<proof>

lemma *indep-events-singleton:*

assumes *F A \in events*

shows *indep-events F {A}*

<proof>

lemma *indep-events-ev-pairs:*

assumes *indep-events F S*

assumes *A \in S B \in S A \neq B*

shows *indep-event (F A) (F B)*

<proof>

lemma *indep-events-ev-inter-pairs:*

assumes *indep-events F S*

assumes *finite A finite B*

assumes *A \neq {} B \neq {}*

assumes $A \subseteq S \ B \subseteq S \ A \cap B = \{\}$
shows $\text{indep-event } (\bigcap (F \text{ ' } A)) (\bigcap (F \text{ ' } B))$
 <proof>

lemma *indep-events-ev-inter-single*:
assumes *indep-events* $F \ S$
assumes *finite* B
assumes $B \neq \{\}$
assumes $A \in S \ B \subseteq S \ A \notin B$
shows $\text{indep-event } (F \ A) (\bigcap (F \text{ ' } B))$
 <proof>

lemma *indep-events-fn-eq*:
assumes $\bigwedge Ai. Ai \in E \implies F \ Ai = G \ Ai$
assumes *indep-events* $F \ E$
shows *indep-events* $G \ E$
 <proof>

lemma *indep-events-fn-eq-iff*:
assumes $\bigwedge Ai. Ai \in E \implies F \ Ai = G \ Ai$
shows $\text{indep-events } F \ E \longleftrightarrow \text{indep-events } G \ E$
 <proof>

lemma *indep-events-one-compl*:
assumes *indep-events* $F \ S$
assumes $A \in S$
shows $\text{indep-events } (\lambda i. \text{if } (i = A) \text{ then } (\text{space } M - F \ i) \text{ else } F \ i) \ S$ (**is**
indep-events $?G \ S$)
 <proof>

lemma *indep-events-update-compl*:
assumes *indep-events* $F \ E$
assumes $E = A \cup B$
assumes $A \cap B = \{\}$
assumes *finite* E
shows $\text{indep-events } (\lambda Ai. \text{if } (Ai \in A) \text{ then } (\text{space } M - (F \ Ai)) \text{ else } (F \ Ai)) \ E$
 <proof>

lemma *indep-events-compl*:
assumes *indep-events* $F \ E$
assumes *finite* E
shows $\text{indep-events } (\lambda Ai. \text{space } M - F \ Ai) \ E$
 <proof>

lemma *indep-events-impl-inj-on*:
assumes *finite* A
assumes *indep-events* $F \ A$
assumes $\bigwedge A'. A' \in A \implies \text{prob } (F \ A') > 0 \wedge \text{prob } (F \ A') < 1$

shows *inj-on* $F A$
 ⟨*proof*⟩

lemma *indep-events-imp-set*:
assumes *finite* A
assumes *indep-events* $F A$
assumes $\bigwedge A' . A' \in A \implies \text{prob } (F A') > 0 \wedge \text{prob } (F A') < 1$
shows *indep-events-set* $(F ' A)$
 ⟨*proof*⟩

lemma *indep-event-set-equiv-bij*:
assumes *bij-betw* $F A E$
assumes *finite* E
shows *indep-events-set* $E \longleftrightarrow \text{indep-events } F A$
 ⟨*proof*⟩

5.3 Mutual Independent Events

Note, set based version only if no duplicates in usage case. The `mutual_indep_events` definition is more general and recommended

definition *mutual-indep-set*:: 'a set \implies 'a set set \implies bool
where *mutual-indep-set* $A S \longleftrightarrow A \in \text{events} \wedge S \subseteq \text{events} \wedge (\forall T \subseteq S . T \neq \{\} \implies \text{prob } (A \cap (\bigcap T)) = \text{prob } A * \text{prob } (\bigcap T))$

lemma *mutual-indep-setI[intro]*: $A \in \text{events} \implies S \subseteq \text{events} \implies (\bigwedge T . T \subseteq S \implies T \neq \{\} \implies \text{prob } (A \cap (\bigcap T)) = \text{prob } A * \text{prob } (\bigcap T)) \implies \text{mutual-indep-set } A S$
 ⟨*proof*⟩

lemma *mutual-indep-setD[dest]*: $\text{mutual-indep-set } A S \implies T \subseteq S \implies T \neq \{\} \implies \text{prob } (A \cap (\bigcap T)) = \text{prob } A * \text{prob } (\bigcap T)$
 ⟨*proof*⟩

lemma *mutual-indep-setD2[dest]*: $\text{mutual-indep-set } A S \implies A \in \text{events}$
 ⟨*proof*⟩

lemma *mutual-indep-setD3[dest]*: $\text{mutual-indep-set } A S \implies S \subseteq \text{events}$
 ⟨*proof*⟩

lemma *mutual-indep-subset*: $\text{mutual-indep-set } A S \implies T \subseteq S \implies \text{mutual-indep-set } A T$
 ⟨*proof*⟩

lemma *mutual-indep-event-set-defD*:
assumes *mutual-indep-set* $A S$
assumes *finite* T
assumes $T \subseteq S$

assumes $T \neq \{\}$
shows *indep-event* $A (\bigcap T)$
 ⟨*proof*⟩

lemma *mutual-indep-event-defI*: $A \in \text{events} \implies S \subseteq \text{events} \implies (\bigwedge T. T \subseteq S \implies T \neq \{\}) \implies$
 $\text{indep-event } A (\bigcap T) \implies \text{mutual-indep-set } A S$
 ⟨*proof*⟩

lemma *mutual-indep-singleton-event*: $\text{mutual-indep-set } A S \implies B \in S \implies \text{indep-event } A B$
 ⟨*proof*⟩

lemma *mutual-indep-cond*:
assumes $A \in \text{events}$ **and** $T \subseteq \text{events}$ **and** *finite* T
and *mutual-indep-set* $A S$ **and** $T \subseteq S$ **and** $T \neq \{\}$ **and** $\text{prob } (\bigcap T) \neq 0$
shows $\mathcal{P}(A | (\bigcap T)) = \text{prob } A$
 ⟨*proof*⟩

lemma *mutual-indep-cond-full*:
assumes $A \in \text{events}$ **and** $S \subseteq \text{events}$ **and** *finite* S
and *mutual-indep-set* $A S$ **and** $S \neq \{\}$ **and** $\text{prob } (\bigcap S) \neq 0$
shows $\mathcal{P}(A | (\bigcap S)) = \text{prob } A$
 ⟨*proof*⟩

lemma *mutual-indep-cond-single*:
assumes $A \in \text{events}$ **and** $B \in \text{events}$
and *mutual-indep-set* $A S$ **and** $B \in S$ **and** $\text{prob } B \neq 0$
shows $\mathcal{P}(A | B) = \text{prob } A$
 ⟨*proof*⟩

lemma *mutual-indep-set-empty*: $A \in \text{events} \implies \text{mutual-indep-set } A \{\}$
 ⟨*proof*⟩

lemma *not-mutual-indep-set-itself*:
assumes $\text{prob } A > 0$ **and** $\text{prob } A < 1$
shows $\neg \text{mutual-indep-set } A \{A\}$
 ⟨*proof*⟩

lemma *is-mutual-indep-set-itself*:
assumes $A \in \text{events}$
assumes $\text{prob } A = 0 \vee \text{prob } A = 1$
shows *mutual-indep-set* $A \{A\}$
 ⟨*proof*⟩

lemma *mutual-indep-set-singleton*:
assumes *indep-event* $A B$
shows *mutual-indep-set* $A \{B\}$
 ⟨*proof*⟩

lemma *mutual-indep-set-one-compl*:
assumes *mutual-indep-set* A S
assumes *finite* S
assumes $B \in S$
shows *mutual-indep-set* A $(\{space\ M - B\} \cup S)$
 $\langle proof \rangle$

lemma *mutual-indep-events-set-update-compl*:
assumes *mutual-indep-set* X E
assumes $E = A \cup B$
assumes $A \cap B = \{\}$
assumes *finite* E
shows *mutual-indep-set* X $(((-) (space\ M) ' A) \cup B)$
 $\langle proof \rangle$

lemma *mutual-indep-events-compl*:
assumes *finite* S
assumes *mutual-indep-set* A S
shows *mutual-indep-set* A $((\lambda\ s.\ space\ M - s) ' S)$
 $\langle proof \rangle$

lemma *mutual-indep-set-all*:
assumes $A \subseteq events$
assumes $\bigwedge Ai. Ai \in A \implies (mutual-indep-set\ Ai\ (A - \{Ai\}))$
shows *indep-events-set* A
 $\langle proof \rangle$

Prefered version using indexed notation

definition *mutual-indep-events*:: 'a set \implies (nat \implies 'a set) \implies nat set \implies bool
where *mutual-indep-events* $A\ F\ I \iff A \in events \wedge (F ' I \subseteq events) \wedge (\forall J \subseteq I. J \neq \{\} \implies prob\ (A \cap (\bigcap j \in J. F\ j)) = prob\ A * prob\ (\bigcap j \in J. F\ j))$

lemma *mutual-indep-eventsI[intro]*: $A \in events \implies (F ' I \subseteq events) \implies (\bigwedge J. J \subseteq I \implies J \neq \{\} \implies prob\ (A \cap (\bigcap j \in J. F\ j)) = prob\ A * prob\ (\bigcap j \in J. F\ j)) \implies mutual-indep-events\ A\ F\ I$
 $\langle proof \rangle$

lemma *mutual-indep-eventsD[dest]*: $mutual-indep-events\ A\ F\ I \implies J \subseteq I \implies J \neq \{\} \implies prob\ (A \cap (\bigcap j \in J. F\ j)) = prob\ A * prob\ (\bigcap j \in J. F\ j)$
 $\langle proof \rangle$

lemma *mutual-indep-eventsD2[dest]*: $mutual-indep-events\ A\ F\ I \implies A \in events$
 $\langle proof \rangle$

lemma *mutual-indep-eventsD3[dest]*: $mutual-indep-events\ A\ F\ I \implies F ' I \subseteq$

events
<proof>

lemma *mutual-indep-ev-subset*: *mutual-indep-events* $A \ F \ I \implies J \subseteq I \implies$ *mutual-indep-events* $A \ F \ J$
<proof>

lemma *mutual-indep-event-defD*:
assumes *mutual-indep-events* $A \ F \ I$
assumes *finite* J
assumes $J \subseteq I$
assumes $J \neq \{\}$
shows *indep-event* $A \ (\bigcap j \in J . F \ j)$
<proof>

lemma *mutual-ev-indep-event-defI*: $A \in \text{events} \implies F \ ' \ I \subseteq \text{events} \implies (\bigwedge J . J \subseteq I \implies J \neq \{\}) \implies$
indep-event $A \ (\bigcap (F \ ' \ J)) \implies$ *mutual-indep-events* $A \ F \ I$
<proof>

lemma *mutual-indep-ev-singleton-event*:
assumes *mutual-indep-events* $A \ F \ I$
assumes $B \in F \ ' \ I$
shows *indep-event* $A \ B$
<proof>

lemma *mutual-indep-ev-singleton-event2*:
assumes *mutual-indep-events* $A \ F \ I$
assumes $i \in I$
shows *indep-event* $A \ (F \ i)$
<proof>

lemma *mutual-indep-iff*:
shows *mutual-indep-events* $A \ F \ I \longleftrightarrow$ *mutual-indep-set* $A \ (F \ ' \ I)$
<proof>

lemma *mutual-indep-ev-cond*:
assumes $A \in \text{events}$ **and** $F \ ' \ J \subseteq \text{events}$ **and** *finite* J
and *mutual-indep-events* $A \ F \ I$ **and** $J \subseteq I$ **and** $J \neq \{\}$ **and** $\text{prob} (\bigcap (F \ ' \ J)) \neq 0$
shows $\mathcal{P}(A \ | \ (\bigcap (F \ ' \ J))) = \text{prob } A$
<proof>

lemma *mutual-indep-ev-cond-full*:
assumes $A \in \text{events}$ **and** $F \ ' \ I \subseteq \text{events}$ **and** *finite* I
and *mutual-indep-events* $A \ F \ I$ **and** $I \neq \{\}$ **and** $\text{prob} (\bigcap (F \ ' \ I)) \neq 0$
shows $\mathcal{P}(A \ | \ (\bigcap (F \ ' \ I))) = \text{prob } A$
<proof>

lemma *mutual-indep-ev-cond-single*:

assumes $A \in \text{events}$ **and** $B \in \text{events}$

and *mutual-indep-events* $A \ F \ I$ **and** $B \in F \ ' \ I$ **and** $\text{prob } B \neq 0$

shows $\mathcal{P}(A \ |B) = \text{prob } A$

<proof>

lemma *mutual-indep-ev-empty*: $A \in \text{events} \implies \text{mutual-indep-events } A \ F \ \{\}$

<proof>

lemma *not-mutual-indep-ev-itself*:

assumes $\text{prob } A > 0$ **and** $\text{prob } A < 1$ **and** $A = F \ i$

shows $\neg \text{mutual-indep-events } A \ F \ \{i\}$

<proof>

lemma *is-mutual-indep-ev-itself*:

assumes $A \in \text{events}$ **and** $A = F \ i$

assumes $\text{prob } A = 0 \vee \text{prob } A = 1$

shows *mutual-indep-events* $A \ F \ \{i\}$

<proof>

lemma *mutual-indep-ev-singleton*:

assumes *indep-event* $A \ (F \ i)$

shows *mutual-indep-events* $A \ F \ \{i\}$

<proof>

lemma *mutual-indep-ev-one-compl*:

assumes *mutual-indep-events* $A \ F \ I$

assumes *finite* I

assumes $i \in I$

assumes *space* $M - F \ i = F \ j$

shows *mutual-indep-events* $A \ F \ (\{j\} \cup I)$

<proof>

lemma *mutual-indep-events-update-compl*:

assumes *mutual-indep-events* $X \ F \ S$

assumes $S = A \cup B$

assumes $A \cap B = \{\}$

assumes *finite* S

assumes *bij-betw* $G \ A \ A'$

assumes $\bigwedge i. i \in A \implies F \ (G \ i) = \text{space } M - F \ i$

shows *mutual-indep-events* $X \ F \ (A' \cup B)$

<proof>

lemma *mutual-indep-ev-events-compl*:

assumes *finite* S

assumes *mutual-indep-events* $A \ F \ S$

assumes *bij-betw* $G \ S \ S'$

assumes $\bigwedge i. i \in S \implies F \ (G \ i) = \text{space } M - F \ i$

shows *mutual-indep-events* $A \ F \ S'$

<proof>

Important lemma on relation between independence and mutual independence of a set

lemma *mutual-indep-ev-set-all*:

assumes $F \text{ ' } I \subseteq \text{events}$

assumes $\bigwedge i. i \in I \implies (\text{mutual-indep-events } (F \ i) \ F \ (I - \{i\}))$

shows *indep-events* $F \ I$

<proof>

end

end

6 The Basic Probabilistic Method Framework

This theory includes all aspects of step (3) and (4) of the basic method framework, which are purely probabilistic

theory *Basic-Method* **imports** *Indep-Events*

begin

6.1 More Set and Multiset lemmas

lemma *card-size-set-mset*: $\text{card } (\text{set-mset } A) \leq \text{size } A$

<proof>

lemma *Union-exists*: $\{a \in A . \exists b \in B . P \ a \ b\} = (\bigcup b \in B . \{a \in A . P \ a \ b\})$

<proof>

lemma *Inter-forall*: $B \neq \{\} \implies \{a \in A . \forall b \in B . P \ a \ b\} = (\bigcap b \in B . \{a \in A . P \ a \ b\})$

<proof>

lemma *function-map-multi-filter-size*:

assumes *image-mset* $F \ (\text{mset-set } A) = B$ **and** *finite* A

shows $\text{card } \{a \in A . P \ (F \ a)\} = \text{size } \{\# \ b \in\# \ B . P \ b \ \#\}$

<proof>

lemma *bij-mset-obtain-set-elem*:

assumes *image-mset* $F \ (\text{mset-set } A) = B$

assumes $b \in\# \ B$

obtains a **where** $a \in A$ **and** $F \ a = b$

<proof>

lemma *bij-mset-obtain-mset-elem*:

assumes *finite* A

assumes *image-mset* $F \ (\text{mset-set } A) = B$

assumes $a \in A$

obtains b **where** $b \in\# \ B$ **and** $F \ a = b$

<proof>

lemma *prod-fn-le1*:

fixes $f :: 'c \Rightarrow ('d :: \{\text{comm-monoid-mult, linordered-semidom}\})$

assumes *finite A*

assumes $A \neq \{\}$

assumes $\bigwedge y. y \in A \implies f y \geq 0 \wedge f y < 1$

shows $(\prod_{x \in A} f x) < 1$

<proof>

context *prob-space*

begin

6.2 Existence Lemmas

lemma *prob-lt-one-obtain*:

assumes $\{e \in \text{space } M . Q e\} \in \text{events}$

assumes $\text{prob } \{e \in \text{space } M . Q e\} < 1$

obtains e **where** $e \in \text{space } M$ **and** $\neg Q e$

<proof>

lemma *prob-gt-zero-obtain*:

assumes $\{e \in \text{space } M . Q e\} \in \text{events}$

assumes $\text{prob } \{e \in \text{space } M . Q e\} > 0$

obtains e **where** $e \in \text{space } M$ **and** $Q e$

<proof>

lemma *inter-gt0-event*:

assumes $F \text{ ' } I \subseteq \text{events}$

assumes $\text{prob } (\bigcap_{i \in I} (\text{space } M - (F i))) > 0$

shows $(\bigcap_{i \in I} (\text{space } M - (F i))) \in \text{events}$ **and** $(\bigcap_{i \in I} (\text{space } M - (F i))) \neq \{\}$

<proof>

lemma *obtain-intersection*:

assumes $F \text{ ' } I \subseteq \text{events}$

assumes $\text{prob } (\bigcap_{i \in I} (\text{space } M - (F i))) > 0$

obtains e **where** $e \in \text{space } M$ **and** $\bigwedge i. i \in I \implies e \notin F i$

<proof>

lemma *obtain-intersection-prop*:

assumes $F \text{ ' } I \subseteq \text{events}$

assumes $\bigwedge i. i \in I \implies F i = \{e \in \text{space } M . P e i\}$

assumes $\text{prob } (\bigcap_{i \in I} (\text{space } M - (F i))) > 0$

obtains e **where** $e \in \text{space } M$ **and** $\bigwedge i. i \in I \implies \neg P e i$

<proof>

lemma *not-in-big-union*:

assumes $\bigwedge i. i \in A \implies e \notin i$

shows $e \notin (\bigcup A)$
<proof>

lemma *not-in-big-union-fn*:
assumes $\bigwedge i . i \in A \implies e \notin F i$
shows $e \notin (\bigcup i \in A . F i)$
<proof>

lemma *obtain-intersection-union*:
assumes $F ' I \subseteq \text{events}$
assumes $\text{prob} (\bigcap i \in I . (\text{space } M - (F i))) > 0$
obtains e **where** $e \in \text{space } M$ **and** $e \notin (\bigcup i \in I . F i)$
<proof>

6.3 Basic Bounds

Lemmas on the Complete Independence and Union bound

lemma *complete-indep-bound1*:
assumes *finite* A
assumes $A \neq \{\}$
assumes $A \subseteq \text{events}$
assumes *indep-events-set* A
assumes $\bigwedge a . a \in A \implies \text{prob } a < 1$
shows $\text{prob} (\text{space } M - (\bigcap A)) > 0$
<proof>

lemma *complete-indep-bound1-index*:
assumes *finite* A
assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *indep-events* $F A$
assumes $\bigwedge a . a \in A \implies \text{prob} (F a) < 1$
shows $\text{prob} (\text{space } M - (\bigcap (F ' A))) > 0$
<proof>

lemma *complete-indep-bound2*:
assumes *finite* A
assumes $A \subseteq \text{events}$
assumes *indep-events-set* A
assumes $\bigwedge a . a \in A \implies \text{prob } a < 1$
shows $\text{prob} (\text{space } M - (\bigcup A)) > 0$
<proof>

lemma *complete-indep-bound2-index*:
assumes *finite* A
assumes $F ' A \subseteq \text{events}$
assumes *indep-events* $F A$
assumes $\bigwedge a . a \in A \implies \text{prob} (F a) < 1$
shows $\text{prob} (\text{space } M - (\bigcup (F ' A))) > 0$

<proof>

lemma *complete-indep-bound3:*

assumes *finite A*
assumes $A \neq \{\}$
assumes $F \text{ ' } A \subseteq \text{events}$
assumes *indep-events F A*
assumes $\bigwedge a . a \in A \implies \text{prob } (F a) < 1$
shows $\text{prob } (\bigcap a \in A. \text{space } M - F a) > 0$
<proof>

Combining complete independence with existence step

lemma *complete-indep-bound-obtain:*

assumes *finite A*
assumes $A \subseteq \text{events}$
assumes *indep-events-set A*
assumes $\bigwedge a . a \in A \implies \text{prob } a < 1$
obtains *e where e ∈ space M and e ∉ ∪ A*
<proof>

lemma *Union-bound-events:*

assumes *finite A*
assumes $A \subseteq \text{events}$
shows $\text{prob } (\bigcup A) \leq (\sum a \in A. \text{prob } a)$
<proof>

lemma *Union-bound-events-fun:*

assumes *finite A*
assumes $f \text{ ' } A \subseteq \text{events}$
shows $\text{prob } (\bigcup (f \text{ ' } A)) \leq (\sum a \in A. \text{prob } (f a))$
<proof>

lemma *Union-bound-avoid:*

assumes *finite A*
assumes $(\sum a \in A. \text{prob } a) < 1$
assumes $A \subseteq \text{events}$
shows $\text{prob } (\text{space } M - \bigcup A) > 0$
<proof>

lemma *Union-bound-avoid-fun:*

assumes *finite A*
assumes $(\sum a \in A. \text{prob } (f a)) < 1$
assumes $f \text{ ' } A \subseteq \text{events}$
shows $\text{prob } (\text{space } M - \bigcup (f \text{ ' } A)) > 0$
<proof>

Combining union bound with existence step

lemma *Union-bound-obtain:*

```

assumes finite A
assumes  $(\sum a \in A. \text{prob } a) < 1$ 
assumes  $A \subseteq \text{events}$ 
obtains e where  $e \in \text{space } M$  and  $e \notin \bigcup A$ 
<proof>

```

```

lemma Union-bound-obtain-fun:
assumes finite A
assumes  $(\sum a \in A. \text{prob } (f \ a)) < 1$ 
assumes  $f' \ A \subseteq \text{events}$ 
obtains e where  $e \in \text{space } M$  and  $e \notin \bigcup (f' \ A)$ 
<proof>

```

```

lemma Union-bound-obtain-compl:
assumes finite A
assumes  $(\sum a \in A. \text{prob } a) < 1$ 
assumes  $A \subseteq \text{events}$ 
obtains e where  $e \in (\text{space } M - \bigcup A)$ 
<proof>

```

```

lemma Union-bound-obtain-compl-fun:
assumes finite A
assumes  $(\sum a \in A. \text{prob } (f \ a)) < 1$ 
assumes  $f' \ A \subseteq \text{events}$ 
obtains e where  $e \in (\text{space } M - \bigcup (f' \ A))$ 
<proof>

```

end

end

7 Lovasz Local Lemma

```

theory Lovasz-Local-Lemma
imports
  Basic-Method
  HOL-Real-Asymp.Real-Asymp
  Indep-Events
  Digraph-Extensions
begin

```

7.1 Random Lemmas on Product Operator

```

lemma prod-constant-ge:
fixes  $y :: 'b :: \{\text{comm-monoid-mult}, \text{linordered-semidom}\}$ 
assumes  $\text{card } A \leq k$ 
assumes  $y \geq 0$  and  $y < 1$ 
shows  $(\prod x \in A. y) \geq y \wedge k$ 
<proof>

```

lemma (in *linordered-idom*) *prod-mono3*:
assumes *finite J I* $I \subseteq J \wedge i. i \in J \implies 0 \leq f i$ ($\wedge i. i \in J \implies f i \leq 1$)
shows $\text{prod } f J \leq \text{prod } f I$
 $\langle \text{proof} \rangle$

lemma *bij-on-ss-image*:
assumes $A \subseteq B$
assumes *bij-betw g B B'*
shows $g ` A \subseteq B'$
 $\langle \text{proof} \rangle$

lemma *bij-on-ss-proper-image*:
assumes $A \subset B$
assumes *bij-betw g B B'*
shows $g ` A \subset B'$
 $\langle \text{proof} \rangle$

7.2 Dependency Graph Concept

Uses directed graphs. The *pair_digraph* locale was sufficient as multi-edges are irrelevant

locale *dependency-digraph* = *pair-digraph G :: nat pair-pre-digraph + prob-space M :: 'a measure*
for $G M$ + **fixes** $F :: \text{nat} \Rightarrow \text{'a set}$
assumes *vss*: $F ` (\text{pverts } G) \subseteq \text{events}$
assumes *mis*: $\wedge i. i \in (\text{pverts } G) \implies \text{mutual-indep-events } (F i) F ((\text{pverts } G) - (\{i\} \cup \text{neighborhood } i))$
begin

lemma *dep-graph-indiv-nh-indep*:
assumes $A \in \text{pverts } G$ $B \in \text{pverts } G$
assumes $B \notin \text{neighborhood } A$
assumes $A \neq B$
assumes $\text{prob } (F B) \neq 0$
shows $\mathcal{P}((F A) \mid (F B)) = \text{prob } (F A)$
 $\langle \text{proof} \rangle$

lemma *mis-subset*:
assumes $i \in \text{pverts } G$
assumes $A \subseteq \text{pverts } G$
shows $\text{mutual-indep-events } (F i) F (A - (\{i\} \cup \text{neighborhood } i))$
 $\langle \text{proof} \rangle$

lemma *dep-graph-indep-events*:
assumes $A \subseteq \text{pverts } G$
assumes $\wedge Ai. Ai \in A \implies \text{out-degree } G Ai = 0$
shows $\text{indep-events } F A$
 $\langle \text{proof} \rangle$

end

7.3 Lovasz Local General Lemma

context *prob-space*

begin

lemma *compl-sets-index*:

assumes $F \text{ ' } A \subseteq \text{events}$

shows $(\lambda i. \text{space } M - F i) \text{ ' } A \subseteq \text{events}$

$\langle \text{proof} \rangle$

lemma *lovasz-inductive-base*:

assumes *dependency-digraph* $G M F$

assumes $\bigwedge Ai. Ai \in A \implies g Ai \geq 0 \wedge g Ai < 1$

assumes $\bigwedge Ai. Ai \in A \implies (\text{prob } (F Ai) \leq (g Ai) * (\prod Aj \in \text{pre-digraph.neighborhood } G Ai. (1 - (g Aj))))$

assumes $Ai \in A$

assumes *pverts* $G = A$

shows $\text{prob } (F Ai) \leq g Ai$

$\langle \text{proof} \rangle$

lemma *lovasz-inductive-base-set*:

assumes $N \subseteq A$

assumes $\bigwedge Ai. Ai \in A \implies g Ai \geq 0 \wedge g Ai < 1$

assumes $\bigwedge Ai. Ai \in A \implies (\text{prob } (F Ai) \leq (g Ai) * (\prod Aj \in N. (1 - (g Aj))))$

assumes $Ai \in A$

shows $\text{prob } (F Ai) \leq g Ai$

$\langle \text{proof} \rangle$

lemma *split-prob-lt-helper*:

assumes *dep-graph*: *dependency-digraph* $G M F$

assumes *dep-graph-verts*: *pverts* $G = A$

assumes *fbounds*: $\bigwedge i. i \in A \implies f i \geq 0 \wedge f i < 1$

assumes *prob-Ai*: $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq$

$(f Ai) * (\prod Aj \in \text{pre-digraph.neighborhood } G Ai. (1 - (f Aj)))$

assumes *aiin*: $Ai \in A$

assumes $N \subseteq \text{pre-digraph.neighborhood } G Ai$

assumes $\exists P1 P2. \mathcal{P}(F Ai \mid \bigcap Aj \in S. \text{space } M - F Aj) = P1/P2 \wedge$

$P1 \leq \text{prob } (F Ai) \wedge P2 \geq (\prod Aj \in N. (1 - (f Aj)))$

shows $\mathcal{P}(F Ai \mid \bigcap Aj \in S. \text{space } M - F Aj) \leq f Ai$

$\langle \text{proof} \rangle$

lemma *lovasz-inequality*:

assumes *finS*: *finite* S

assumes *sevents*: $F \text{ ' } S \subseteq \text{events}$

assumes *S-subset*: $S \subseteq A - \{Ai\}$

assumes *prob2*: $\text{prob } (\bigcap Aj \in S. (\text{space } M - (F Aj))) > 0$

assumes *irange*: $i \in \{0..<\text{card } S1\}$
assumes *bb*: *bij-betw* $g \{0..<\text{card } S1\} S1$
assumes *s1-def*: $S1 = (S \cap N)$
assumes *s2-def*: $S2 = S - S1$
assumes *ne-cond*: $i > 0 \vee S2 \neq \{\}$
assumes *hyps*: $\bigwedge B. B \subset S \implies g \ i \in A \implies B \subseteq A - \{g \ i\} \implies B \neq \{\} \implies$
 $0 < \text{prob} (\bigcap_{Aj \in B. \text{space } M - F \ Aj}) \implies \mathcal{P}(F \ (g \ i) \mid \bigcap_{Aj \in B. \text{space } M - F$
 $Aj) \leq f \ (g \ i)$
shows $\mathcal{P}((\text{space } M - F \ (g \ i)) \mid (\bigcap ((\lambda \ i. \text{space } M - F \ i) \ 'g \ '\{0..<i\} \cup ((\lambda \ i.$
 $\text{space } M - F \ i) \ 'S2))))$
 $\geq (1 - f \ (g \ i))$
<proof>

The main helper lemma

lemma *lovasz-inductive*:

assumes *finA*: *finite* A
assumes *Aevents*: $F \ 'A \subseteq \text{events}$
assumes *fbounds*: $\bigwedge i. i \in A \implies f \ i \geq 0 \wedge f \ i < 1$
assumes *dep-graph*: *dependency-digraph* $G \ M \ F$
assumes *dep-graph-verts*: *pverts* $G = A$
assumes *prob-Ai*: $\bigwedge Ai. Ai \in A \implies \text{prob} (F \ Ai) \leq$
 $(f \ Ai) * (\prod_{Aj \in \text{pre-digraph.neighborhood } G \ Ai. (1 - (f \ Aj)))$
assumes *Ai-in*: $Ai \in A$
assumes *S-subset*: $S \subseteq A - \{Ai\}$
assumes *S-nempty*: $S \neq \{\}$
assumes *prob2*: $\text{prob} (\bigcap_{Aj \in S. (\text{space } M - (F \ Aj)))} > 0$
shows $\mathcal{P}((F \ Ai) \mid (\bigcap_{Aj \in S. (\text{space } M - (F \ Aj)))) \leq f \ Ai$
<proof>

The main lemma

theorem *lovasz-local-general*:

assumes $A \neq \{\}$
assumes *F ' A* $\subseteq \text{events}$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies f \ Ai \geq 0 \wedge f \ Ai < 1$
assumes *dependency-digraph* $G \ M \ F$
assumes $\bigwedge Ai. Ai \in A \implies (\text{prob} (F \ Ai) \leq (f \ Ai) * (\prod_{Aj \in \text{pre-digraph.neighborhood}$
 $G \ Ai. (1 - (f \ Aj))))$
assumes *pverts* $G = A$
shows $\text{prob} (\bigcap_{Ai \in A. (\text{space } M - (F \ Ai))) \geq (\prod_{Ai \in A. (1 - f \ Ai)}) (\prod_{Ai \in A. (1 - f \ Ai)} > 0$
<proof>

7.4 Lovasz Corollaries and Variations

corollary *lovasz-local-general-positive*:

assumes $A \neq \{\}$
assumes *F ' A* $\subseteq \text{events}$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies f \ Ai \geq 0 \wedge f \ Ai < 1$

assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies (\text{prob } (F Ai) \leq$
 $(f Ai) * (\prod Aj \in \text{pre-digraph.neighborhood } G Ai. (1 - (f Aj))))$
assumes *pverts* $G = A$
shows $\text{prob } (\bigcap Ai \in A . (\text{space } M - (F Ai))) > 0$
 $\langle \text{proof} \rangle$

theorem *lovasz-local-symmetric-dep-graph*:

fixes $e :: \text{real}$
fixes $d :: \text{nat}$
assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *finite* A
assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies \text{out-degree } G Ai \leq d$
assumes $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$
assumes $\exp(1) * p * (d + 1) \leq 1$
assumes *pverts* $G = A$
shows $\text{prob } (\bigcap Ai \in A . (\text{space } M - (F Ai))) > 0$
 $\langle \text{proof} \rangle$

corollary *lovasz-local-symmetric4gt*:

fixes $e :: \text{real}$
fixes $d :: \text{nat}$
assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *finite* A
assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies \text{out-degree } G Ai \leq d$
assumes $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$
assumes $\frac{1}{4} * p * d \leq 1$
assumes $d \geq 3$
assumes *pverts* $G = A$
shows $\text{prob } (\bigcap Ai \in A . (\text{space } M - F Ai)) > 0$
 $\langle \text{proof} \rangle$

lemma *lovasz-local-symmetric4*:

fixes $e :: \text{real}$
fixes $d :: \text{nat}$
assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *finite* A
assumes *dependency-digraph* $G M F$
assumes $\bigwedge Ai. Ai \in A \implies \text{out-degree } G Ai \leq d$
assumes $\bigwedge Ai. Ai \in A \implies \text{prob } (F Ai) \leq p$
assumes $\frac{1}{4} * p * d \leq 1$
assumes $d \geq 1$
assumes *pverts* $G = A$

shows $\text{prob} (\bigcap Ai \in A . (\text{space } M - F Ai)) > 0$
 ⟨proof⟩

Converting between dependency graph and indexed set representation of mutual independence

lemma (in *pair-digraph*) *g-Ai-simplification*:

assumes $Ai \in A$
assumes $g Ai \subseteq A - \{Ai\}$
assumes $pverts G = A$
assumes $\text{parcs } G = \{e \in A \times A . \text{snd } e \in (A - (\{fst e\} \cup (g (fst e))))\}$
shows $g Ai = A - (\{Ai\} \cup \text{neighborhood } Ai)$
 ⟨proof⟩

lemma *define-dep-graph-set*:

assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies g Ai \subseteq A - \{Ai\} \wedge \text{mutual-indep-events } (F Ai) F (g Ai)$
shows $\text{dependency-digraph } (\text{pverts} = A, \text{parcs} = \{e \in A \times A . \text{snd } e \in (A - (\{fst e\} \cup (g (fst e))))\}) \text{ } M F$
 (is *dependency-digraph* ? $G M F$)
 ⟨proof⟩

lemma *define-dep-graph-deg-bound*:

assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies g Ai \subseteq A - \{Ai\} \wedge \text{card } (g Ai) \geq \text{card } A - d - 1$
 \wedge
 $\text{mutual-indep-events } (F Ai) F (g Ai)$
shows $\bigwedge Ai. Ai \in A \implies$
 $\text{out-degree } (\text{pverts} = A, \text{parcs} = \{e \in A \times A . \text{snd } e \in (A - (\{fst e\} \cup (g (fst e))))\}) \text{ } Ai \leq d$
 (is $\bigwedge Ai. Ai \in A \implies \text{out-degree } (\text{with-proj } ?G) Ai \leq d$)
 ⟨proof⟩

lemma *obtain-dependency-graph*:

assumes $A \neq \{\}$
assumes $F ' A \subseteq \text{events}$
assumes *finite* A
assumes $\bigwedge Ai. Ai \in A \implies$
 $(\exists S . S \subseteq A - \{Ai\} \wedge \text{card } S \geq \text{card } A - d - 1 \wedge \text{mutual-indep-events } (F Ai) F S)$
obtains G where $\text{dependency-digraph } G M F \text{ pverts } G = A \wedge \bigwedge Ai. Ai \in A \implies$
 $\text{out-degree } G Ai \leq d$
 ⟨proof⟩

This is the variation of the symmetric version most commonly in use

theorem *lovasz-local-symmetric*:

```

fixes  $d :: nat$ 
assumes  $A \neq \{\}$ 
assumes  $F \text{ ' } A \subseteq events$ 
assumes finite A
assumes  $\bigwedge Ai. Ai \in A \implies (\exists S . S \subseteq A - \{Ai\} \wedge card S \geq card A - d - 1$ 
 $\wedge mutual-indep-events (F Ai) F S)$ 
assumes  $\bigwedge Ai. Ai \in A \implies prob (F Ai) \leq p$ 
assumes  $exp(1) * p * (d + 1) \leq 1$ 
shows  $prob (\bigcap Ai \in A . (space M - (F Ai))) > 0$ 
<proof>

```

lemma *lovasz-local-symmetric4-set*:

```

fixes  $d :: nat$ 
assumes  $A \neq \{\}$ 
assumes  $F \text{ ' } A \subseteq events$ 
assumes finite A
assumes  $\bigwedge Ai. Ai \in A \implies (\exists S . S \subseteq A - \{Ai\} \wedge card S \geq card A - d - 1$ 
 $\wedge mutual-indep-events (F Ai) F S)$ 
assumes  $\bigwedge Ai. Ai \in A \implies prob (F Ai) \leq p$ 
assumes  $\frac{1}{4} * p * d \leq 1$ 
assumes  $d \geq 1$ 
shows  $prob (\bigcap Ai \in A . (space M - F Ai)) > 0$ 
<proof>
end

```

end

theory *Lovasz-Local-Root*

imports

PiE-Rel-Extras

Digraph-Extensions

Prob-Events-Extras

Cond-Prob-Extensions

Indep-Events

Basic-Method

Lovasz-Local-Lemma

begin

end

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