

Linear Inequalities*

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Abstract

We formalize results about linear inequalities, mainly from Schrijver's book [3]. The main results are the proof of the fundamental theorem on linear inequalities, Farkas' lemma, Carathéodory's theorem, the Farkas-Minkowsky-Weyl theorem, the decomposition theorem of polyhedra, and Meyer's result that the integer hull of a polyhedron is a polyhedron itself. Several theorems include bounds on the appearing numbers, and in particular we provide an a-priori bound on mixed-integer solutions of linear inequalities.

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1 Introduction

The motivation for this formalization is the aim of developing a verified theory solver for linear integer arithmetic. Such a solver can be a combination of a simplex-implementation within a branch-and-bound approach, that might also utilize Gomory cuts [1, Section 4 of the extended version]. However, the branch-and-bound algorithm does not terminate in general, since the search space is infinite. To solve this latter problem, one can use results of Papadimitriou: he showed that whenever a set of linear inequalities has an integer solution, then it also has a small solution, where the bound on such a solution can be computed easily from the input [2].

In this entry, we therefore formalize several results on linear inequalities which are required to obtain the desired bound, by following the proofs of Schrijver's textbook [3, Sections 7 and 16].

We start with basic definitions and results on cones, convex hulls, and polyhedra. Next, we verify the fundamental theorem of linear inequalities, which in our formalization shows the equivalence of four statements to describe a cone. From this theorem, one easily derives Farkas' Lemma and Carathéodory's theorem. Moreover we verify the Farkas-Minkowsky-Weyl theorem, that a convex cone is polyhedral if and only if it is finitely generated, and use this result to obtain the decomposition theorem for polyhedra, i.e., that a polyhedron can always be decomposed into a polytope and a finitely generated cone. For most of the previously mentioned results, we include bounds, so that in particular we have a quantitative version of the decomposition theorem, which provides bounds on the vectors that construct the polytope and the cone, and where these bounds are computed directly from the input polyhedron that should be decomposed.

We further prove the decomposition theorem also for the integer hull of a polyhedron, using the same bounds, which gives rise to small integer solutions for linear inequalities. We finally formalize a direct proof for the

more general case of mixed integer solutions, where we also permit both strict and non-strict linear inequalities.

Theorem 1. Consider $A_1 \in \mathbb{Z}^{m_1 \times n}$, $b_1 \in \mathbb{Z}^{m_1}$, $A_2 \in \mathbb{Z}^{m_2 \times n}$, $b_2 \in \mathbb{Z}^{m_2}$. Let β be a bound on A_1, b_1, A_2, b_2 , i.e., $\beta \geq |z|$ for all numbers z that occur within A_1, b_1, A_2, b_2 . Let $n = n_1 + n_2$. Then if $x \in \mathbb{Z}^{n_1} \times \mathbb{R}^{n_2} \subseteq \mathbb{R}^n$ is a mixed integer solution of the linear inequalities, i.e., $A_1 x \leq b_1$ and $A_2 x < b_2$, then there also exists a mixed integer solution $y \in \mathbb{Z}^{n_1} \times \mathbb{R}^{n_2}$ where $|y_i| \leq (n+1)! \cdot \beta^n$ for each entry y_i of y .

The verified bound in Theorem 1 in particular implies that integer-satisfiability of linear-inequalities with integer coefficients is in NP.

2 Missing Lemmas on Vectors and Matrices

We provide some results on vector spaces which should be merged into Jordan-Normal-Form/Matrix.

theory *Missing-Matrix*

imports *Jordan-Normal-Form.Matrix*

begin

lemma *orthogonalD'*: **assumes** *orthogonal vs*

and $v \in \text{set } vs$ **and** $w \in \text{set } vs$

shows $(v \cdot w = 0) = (v \neq w)$

<proof>

lemma *zero-mat-mult-vector[simp]*: $x \in \text{carrier-vec } nc \implies 0_m \text{ nr } nc *_{\cdot} x = 0_v$
nr

<proof>

lemma *add-diff-cancel-right-vec*:

$a \in \text{carrier-vec } n \implies (b :: 'a :: \text{cancel-ab-semigroup-add vec}) \in \text{carrier-vec } n \implies$

$(a + b) - b = a$

<proof>

lemma *elements-four-block-mat-id*:

assumes $c: A \in \text{carrier-mat } nr1 \ nc1$ $B \in \text{carrier-mat } nr1 \ nc2$

$C \in \text{carrier-mat } nr2 \ nc1$ $D \in \text{carrier-mat } nr2 \ nc2$

shows

$\text{elements-mat } (\text{four-block-mat } A \ B \ C \ D) =$

$\text{elements-mat } A \cup \text{elements-mat } B \cup \text{elements-mat } C \cup \text{elements-mat } D$

(**is** $\text{elements-mat } ?four = ?X$)

<proof>

lemma *elements-mat-append-rows*: $A \in \text{carrier-mat } nr \ n \implies B \in \text{carrier-mat } nr^2$
 $n \implies$

$\text{elements-mat } (A @_r B) = \text{elements-mat } A \cup \text{elements-mat } B$

<proof>

lemma *elements-mat-uminus[simp]*: *elements-mat* $(-A) = \text{uminus } 'elements-mat$
A

<proof>

lemma *vec-set-uminus[simp]*: *vec-set* $(-A) = \text{uminus } 'vec-set$ *A*

<proof>

definition *append-cols* :: *'a* :: *zero mat* \Rightarrow *'a mat* \Rightarrow *'a mat* (**infixr** $@_c$ 65) **where**
 $A @_c B = (A^T @_r B^T)^T$

lemma *carrier-append-cols[simp, intro]*:

$A \in \text{carrier-mat } nr \ nc1 \Longrightarrow$

$B \in \text{carrier-mat } nr \ nc2 \Longrightarrow (A @_c B) \in \text{carrier-mat } nr \ (nc1 + nc2)$

<proof>

lemma *elements-mat-transpose-mat[simp]*: *elements-mat* $(A^T) = \text{elements-mat } A$

<proof>

lemma *elements-mat-append-cols*: $A \in \text{carrier-mat } n \ nc \Longrightarrow B \in \text{carrier-mat } n$
nc1

$\Longrightarrow \text{elements-mat } (A @_c B) = \text{elements-mat } A \cup \text{elements-mat } B$

<proof>

lemma *vec-first-index*:

assumes $v: \text{dim-vec } v \geq n$

and $i: i < n$

shows $(\text{vec-first } v \ n) \$ i = v \$ i$

<proof>

lemma *vec-last-index*:

assumes $v: v \in \text{carrier-vec } (n + m)$

and $i: i < m$

shows $(\text{vec-last } v \ m) \$ i = v \$ (n + i)$

<proof>

lemma *vec-first-add*:

assumes $\text{dim-vec } x \geq n$

and $\text{dim-vec } y \geq n$

shows $\text{vec-first } (x + y) \ n = \text{vec-first } x \ n + \text{vec-first } y \ n$

<proof>

lemma *vec-first-zero[simp]*: $m \leq n \Longrightarrow \text{vec-first } (0_v \ n) \ m = 0_v \ m$

<proof>

lemma *vec-first-smult*:

$\llbracket m \leq n; x \in \text{carrier-vec } n \rrbracket \Longrightarrow \text{vec-first } (c \cdot_v x) \ m = c \cdot_v \text{vec-first } x \ m$

<proof>

lemma *elements-mat-mat-of-row*[simp]: $\text{elements-mat } (\text{mat-of-row } v) = \text{vec-set } v$
 ⟨proof⟩

lemma *vec-set-append-vec*[simp]: $\text{vec-set } (v @_v w) = \text{vec-set } v \cup \text{vec-set } w$
 ⟨proof⟩

lemma *vec-set-vNil*[simp]: $\text{set}_v v\text{Nil} = \{\}$ ⟨proof⟩

lemma *diff-smult-distrib-vec*: $((x :: 'a::\text{ring}) - y) \cdot_v v = x \cdot_v v - y \cdot_v v$
 ⟨proof⟩

lemma *add-diff-eq-vec*: **fixes** $y :: 'a :: \text{group-add } \text{vec}$
shows $y \in \text{carrier-vec } n \implies x \in \text{carrier-vec } n \implies z \in \text{carrier-vec } n \implies y + (x - z) = y + x - z$
 ⟨proof⟩

definition *mat-of-col* $v = (\text{mat-of-row } v)^T$

lemma *elements-mat-mat-of-col*[simp]: $\text{elements-mat } (\text{mat-of-col } v) = \text{vec-set } v$
 ⟨proof⟩

lemma *mat-of-col-dim*[simp]: $\text{dim-row } (\text{mat-of-col } v) = \text{dim-vec } v$
 $\text{dim-col } (\text{mat-of-col } v) = 1$
 $\text{mat-of-col } v \in \text{carrier-mat } (\text{dim-vec } v) 1$
 ⟨proof⟩

lemma *col-mat-of-col*[simp]: $\text{col } (\text{mat-of-col } v) 0 = v$
 ⟨proof⟩

lemma *mult-mat-of-col*: $A \in \text{carrier-mat } nr \ nc \implies v \in \text{carrier-vec } nc \implies$
 $A * \text{mat-of-col } v = \text{mat-of-col } (A *_v v)$
 ⟨proof⟩

lemma *mat-mult-append-cols*: **fixes** $A :: 'a :: \text{comm-semiring-0 } \text{mat}$
assumes $A: A \in \text{carrier-mat } nr \ nc1$
and $B: B \in \text{carrier-mat } nr \ nc2$
and $v1: v1 \in \text{carrier-vec } nc1$
and $v2: v2 \in \text{carrier-vec } nc2$
shows $(A @_c B) *_v (v1 @_v v2) = A *_v v1 + B *_v v2$
 ⟨proof⟩

lemma *vec-first-append*:
assumes $v \in \text{carrier-vec } n$
shows $\text{vec-first } (v @_v w) n = v$
 ⟨proof⟩

lemma *vec-le-iff-diff-le-0*: **fixes** $a :: 'a :: \text{ordered-ab-group-add } \text{vec}$

shows $(a \leq b) = (a - b \leq 0_v \text{ (dim-vec } a))$
 ⟨proof⟩

definition $\text{mat-row-first } A \ n \equiv \text{mat } n \text{ (dim-col } A) \ (\lambda \ (i, j). \ A \ \$\$ \ (i, j))$

definition $\text{mat-row-last } A \ n \equiv \text{mat } n \text{ (dim-col } A) \ (\lambda \ (i, j). \ A \ \$\$ \ (\text{dim-row } A - n + i, j))$

lemma $\text{mat-row-first-carrier[simp]}$: $\text{mat-row-first } A \ n \in \text{carrier-mat } n \text{ (dim-col } A)$
 ⟨proof⟩

lemma $\text{mat-row-first-dim[simp]}$:
 $\text{dim-row } (\text{mat-row-first } A \ n) = n$
 $\text{dim-col } (\text{mat-row-first } A \ n) = \text{dim-col } A$
 ⟨proof⟩

lemma $\text{mat-row-last-carrier[simp]}$: $\text{mat-row-last } A \ n \in \text{carrier-mat } n \text{ (dim-col } A)$
 ⟨proof⟩

lemma $\text{mat-row-last-dim[simp]}$:
 $\text{dim-row } (\text{mat-row-last } A \ n) = n$
 $\text{dim-col } (\text{mat-row-last } A \ n) = \text{dim-col } A$
 ⟨proof⟩

lemma $\text{mat-row-first-nth[simp]}$: $i < n \implies \text{row } (\text{mat-row-first } A \ n) \ i = \text{row } A \ i$
 ⟨proof⟩

lemma append-rows-nth :
assumes $A \in \text{carrier-mat } nr1 \ nc$
and $B \in \text{carrier-mat } nr2 \ nc$
shows $i < nr1 \implies \text{row } (A \ @_r \ B) \ i = \text{row } A \ i$
and $\llbracket i \geq nr1; i < nr1 + nr2 \rrbracket \implies \text{row } (A \ @_r \ B) \ i = \text{row } B \ (i - nr1)$
 ⟨proof⟩

lemma $\text{mat-of-row-last-nth[simp]}$:
 $i < n \implies \text{row } (\text{mat-row-last } A \ n) \ i = \text{row } A \ (\text{dim-row } A - n + i)$
 ⟨proof⟩

lemma $\text{mat-row-first-last-append}$:
assumes $\text{dim-row } A = m + n$
shows $(\text{mat-row-first } A \ m) \ @_r \ (\text{mat-row-last } A \ n) = A$
 ⟨proof⟩

definition $\text{mat-col-first } A \ n \equiv (\text{mat-row-first } A^T \ n)^T$

definition $\text{mat-col-last } A \ n \equiv (\text{mat-row-last } A^T \ n)^T$

lemma $\text{mat-col-first-carrier[simp]}$: $\text{mat-col-first } A \ n \in \text{carrier-mat } (\text{dim-row } A) \ n$
 ⟨proof⟩

lemma *mat-col-first-dim*[simp]:

$$\dim\text{-row } (\text{mat-col-first } A \ n) = \dim\text{-row } A$$

$$\dim\text{-col } (\text{mat-col-first } A \ n) = n$$

<proof>

lemma *mat-col-last-carrier*[simp]: $\text{mat-col-last } A \ n \in \text{carrier-mat } (\dim\text{-row } A) \ n$

<proof>

lemma *mat-col-last-dim*[simp]:

$$\dim\text{-row } (\text{mat-col-last } A \ n) = \dim\text{-row } A$$

$$\dim\text{-col } (\text{mat-col-last } A \ n) = n$$

<proof>

lemma *mat-col-first-nth*[simp]:

$$\llbracket i < n; i < \dim\text{-col } A \rrbracket \implies \text{col } (\text{mat-col-first } A \ n) \ i = \text{col } A \ i$$

<proof>

lemma *append-cols-nth*:

assumes $A \in \text{carrier-mat } nr \ nc1$

and $B \in \text{carrier-mat } nr \ nc2$

shows $i < nc1 \implies \text{col } (A \ @_c \ B) \ i = \text{col } A \ i$

and $\llbracket i \geq nc1; i < nc1 + nc2 \rrbracket \implies \text{col } (A \ @_c \ B) \ i = \text{col } B \ (i - nc1)$

<proof>

lemma *mat-of-col-last-nth*[simp]:

$$\llbracket i < n; i < \dim\text{-col } A \rrbracket \implies \text{col } (\text{mat-col-last } A \ n) \ i = \text{col } A \ (\dim\text{-col } A - n + i)$$

<proof>

lemma *mat-col-first-last-append*:

assumes $\dim\text{-col } A = m + n$

shows $(\text{mat-col-first } A \ m) \ @_c \ (\text{mat-col-last } A \ n) = A$

<proof>

lemma *mat-of-row-dim-row-1*: $(\dim\text{-row } A = 1) = (A = \text{mat-of-row } (\text{row } A \ 0))$

<proof>

lemma *mat-of-col-dim-col-1*: $(\dim\text{-col } A = 1) = (A = \text{mat-of-col } (\text{col } A \ 0))$

<proof>

definition *vec-of-scal* :: $'a \Rightarrow 'a \ \text{vec}$ **where** $\text{vec-of-scal } x \equiv \text{vec } 1 \ (\lambda \ i. \ x)$

lemma *vec-of-scal-dim*[simp]:

$$\dim\text{-vec } (\text{vec-of-scal } x) = 1$$

$$\text{vec-of-scal } x \in \text{carrier-vec } 1$$

<proof>

lemma *index-vec-of-scal*[simp]: $(\text{vec-of-scal } x) \ \$ \ 0 = x$

<proof>

lemma *row-mat-of-col[simp]*: $i < \dim\text{-vec } v \implies \text{row } (\text{mat-of-col } v) \ i = \text{vec-of-scal } (v \ \$ \ i)$
<proof>

lemma *vec-of-scal-dim-1*: $(v \in \text{carrier-vec } 1) = (v = \text{vec-of-scal } (v \ \$ \ 0))$
<proof>

lemma *mult-mat-of-row-vec-of-scal*: **fixes** $x :: 'a :: \text{comm-ring-1}$
shows $\text{mat-of-col } v \ *_v \ \text{vec-of-scal } x = x \cdot_v v$
<proof>

lemma *smult-pos-vec[simp]*: **fixes** $l :: 'a :: \text{linordered-ring-strict}$
assumes $l: l > 0$
shows $(l \cdot_v v \leq 0_v \ n) = (v \leq 0_v \ n)$
<proof>

lemma *finite-elements-mat[simp]*: $\text{finite } (\text{elements-mat } A)$
<proof>

lemma *finite-vec-set[simp]*: $\text{finite } (\text{vec-set } A)$
<proof>

lemma *lesseq-vecI*: **assumes** $v \in \text{carrier-vec } n \ w \in \text{carrier-vec } n$
 $\bigwedge i. i < n \implies v \ \$ \ i \leq w \ \$ \ i$
shows $v \leq w$
<proof>

lemma *lesseq-vecD*: **assumes** $w \in \text{carrier-vec } n$
and $v \leq w$
and $i < n$
shows $v \ \$ \ i \leq w \ \$ \ i$
<proof>

lemma *vec-add-mono*: **fixes** $a :: 'a :: \text{ordered-ab-semigroup-add } \text{vec}$
assumes $\text{dim}: \text{dim-vec } b = \text{dim-vec } d$
and $ab: a \leq b$
and $cd: c \leq d$
shows $a + c \leq b + d$
<proof>

lemma *smult-nneg-npos-vec*: **fixes** $l :: 'a :: \text{ordered-semiring-0}$
assumes $l: l \geq 0$
and $v: v \leq 0_v \ n$
shows $l \cdot_v v \leq 0_v \ n$
<proof>

lemma *smult-vec-nonneg-eq*: **fixes** $c :: 'a :: \text{field}$

shows $c \neq 0 \implies (c \cdot_v x = c \cdot_v y) = (x = y)$
 <proof>

lemma *distinct-smult-nonneg*: **fixes** $c :: 'a :: field$
assumes $c: c \neq 0$
shows $distinct\ lC \implies distinct\ (map\ ((\cdot_v)\ c)\ lC)$
 <proof>

lemma *exists-vec-append*: $(\exists x \in carrier-vec\ (n + m). P\ x) \longleftrightarrow (\exists x1 \in carrier-vec\ n. \exists x2 \in carrier-vec\ m. P\ (x1\ @_v\ x2))$
 <proof>

end

3 Missing Lemmas on Vector Spaces

We provide some results on vector spaces which should be merged into other AFP entries.

theory *Missing-VS-Connect*

imports

Jordan-Normal-Form.VS-Connect

Missing-Matrix

Polynomial-Factorization.Missing-List

begin

context *vec-space*

begin

lemma *span-diff*: **assumes** $A: A \subseteq carrier-vec\ n$

and $a: a \in span\ A$ **and** $b: b \in span\ A$

shows $a - b \in span\ A$

<proof>

lemma *finsum-scalar-prod-sum'*:

assumes $f: f \in U \rightarrow carrier-vec\ n$

and $w: w \in carrier-vec\ n$

shows $w \cdot finsum\ V\ f\ U = sum\ (\lambda u. w \cdot f\ u)\ U$

<proof>

lemma *lincomb-scalar-prod-left*: **assumes** $W \subseteq carrier-vec\ n\ v \in carrier-vec\ n$

shows $lincomb\ a\ W \cdot v = (\sum_{w \in W}. a\ w * (w \cdot v))$

<proof>

lemma *lincomb-scalar-prod-right*: **assumes** $W \subseteq carrier-vec\ n\ v \in carrier-vec\ n$

shows $v \cdot lincomb\ a\ W = (\sum_{w \in W}. a\ w * (v \cdot w))$

<proof>

lemma *lin-indpt-empty[simp]*: $lin-indpt\ \{\}$

<proof>

lemma *span-carrier-lin-indpt-card-n*:
assumes $W \subseteq \text{carrier-vec } n$ $\text{card } W = n$ $\text{lin-indpt } W$
shows $\text{span } W = \text{carrier-vec } n$
 $\langle \text{proof} \rangle$

lemma *ortho-span*: **assumes** $W: W \subseteq \text{carrier-vec } n$
and $X: X \subseteq \text{carrier-vec } n$
and *ortho*: $\bigwedge w x. w \in W \implies x \in X \implies w \cdot x = 0$
and $w: w \in \text{span } W$ **and** $x: x \in X$
shows $w \cdot x = 0$
 $\langle \text{proof} \rangle$

lemma *ortho-span'*: **assumes** $W: W \subseteq \text{carrier-vec } n$
and $X: X \subseteq \text{carrier-vec } n$
and *ortho*: $\bigwedge w x. w \in W \implies x \in X \implies x \cdot w = 0$
and $w: w \in \text{span } W$ **and** $x: x \in X$
shows $x \cdot w = 0$
 $\langle \text{proof} \rangle$

lemma *ortho-span-span*: **assumes** $W: W \subseteq \text{carrier-vec } n$
and $X: X \subseteq \text{carrier-vec } n$
and *ortho*: $\bigwedge w x. w \in W \implies x \in X \implies w \cdot x = 0$
and $w: w \in \text{span } W$ **and** $x: x \in \text{span } X$
shows $w \cdot x = 0$
 $\langle \text{proof} \rangle$

lemma *lincomb-in-span*[*intro*]:
assumes $X: X \subseteq \text{carrier-vec } n$
shows $\text{lincomb } a \ X \in \text{span } X$
 $\langle \text{proof} \rangle$

lemma *generating-card-n-basis*: **assumes** $X: X \subseteq \text{carrier-vec } n$
and *span*: $\text{carrier-vec } n \subseteq \text{span } X$
and *card*: $\text{card } X = n$
shows $\text{basis } X$
 $\langle \text{proof} \rangle$

lemma *lincomb-list-append*:
assumes $Ws: \text{set } Ws \subseteq \text{carrier-vec } n$
shows $\text{set } Vs \subseteq \text{carrier-vec } n \implies \text{lincomb-list } f \ (Vs @ Ws) =$
 $\text{lincomb-list } f \ Vs + \text{lincomb-list } (\lambda i. f \ (i + \text{length } Vs)) \ Ws$
 $\langle \text{proof} \rangle$

lemma *lincomb-list-snoc*[*simp*]:
shows $\text{set } Vs \subseteq \text{carrier-vec } n \implies x \in \text{carrier-vec } n \implies$
 $\text{lincomb-list } f \ (Vs @ [x]) = \text{lincomb-list } f \ Vs + f \ (\text{length } Vs) \cdot_v x$
 $\langle \text{proof} \rangle$

lemma *lincomb-list-smult*:
 $set\ Vs \subseteq carrier\ vec\ n \implies lincomb\ list\ (\lambda\ i.\ a * c\ i)\ Vs = a \cdot_v lincomb\ list\ c\ Vs$
<proof>

lemma *lincomb-list-index*:
assumes $i: i < n$
shows $set\ Xs \subseteq carrier\ vec\ n \implies$
 $lincomb\ list\ c\ Xs\ \$\ i = sum\ (\lambda\ j.\ c\ j * (Xs\ !\ j)\ \$\ i)\ \{0..<length\ Xs\}$
<proof>

end
end

4 Basis Extension

We prove that every linear indepent set/list of vectors can be extended into a basis. Similarly, from every set of vectors one can extract a linear independent set of vectors that spans the same space.

theory *Basis-Extension*
imports
LLL-Basis-Reduction.Gram-Schmidt-2
begin

context *cof-vec-space*
begin

lemma *lin-indpt-list-length-le-n*: **assumes** *lin-indpt-list xs*
shows $length\ xs \leq n$
<proof>

lemma *lin-indpt-list-length-eq-n*: **assumes** *lin-indpt-list xs*
and $length\ xs = n$
shows $span\ (set\ xs) = carrier\ vec\ n\ basis\ (set\ xs)$
<proof>

lemma *expand-to-basis*: **assumes** *lin: lin-indpt-list xs*
shows $\exists\ ys.\ set\ ys \subseteq set\ (unit\ vecs\ n) \wedge lin\ indpt\ list\ (xs\ @\ ys) \wedge length\ (xs\ @\ ys) = n$
<proof>

definition *basis-extension xs = (SOME ys.*
 $set\ ys \subseteq set\ (unit\ vecs\ n) \wedge lin\ indpt\ list\ (xs\ @\ ys) \wedge length\ (xs\ @\ ys) = n$

lemma *basis-extension*: **assumes** *lin-indpt-list xs*
shows $set\ (basis\ extension\ xs) \subseteq set\ (unit\ vecs\ n)$
 $lin\ indpt\ list\ (xs\ @\ basis\ extension\ xs)$
 $length\ (xs\ @\ basis\ extension\ xs) = n$

<proof>

lemma *exists-lin-indpt-sublist*: **assumes** $X \subseteq \text{carrier-vec } n$
shows $\exists Ls. \text{lin-indpt-list } Ls \wedge \text{span } (\text{set } Ls) = \text{span } X \wedge \text{set } Ls \subseteq X$
<proof>

lemma *exists-lin-indpt-subset*: **assumes** $X \subseteq \text{carrier-vec } n$
shows $\exists Ls. \text{lin-indpt } Ls \wedge \text{span } (Ls) = \text{span } X \wedge Ls \subseteq X$
<proof>
end

end

5 Sum of Vector Sets

We use Isabelle's Set-Algebra theory to be able to write $V + W$ for sets of vectors V and W , and prove some obvious properties about them.

theory *Sum-Vec-Set*
imports
 Missing-Matrix
 HOL-Library.Set-Algebras
begin

lemma *add-0-right-vecset*:
assumes $(A :: 'a :: \text{monoid-add vec set}) \subseteq \text{carrier-vec } n$
shows $A + \{0_v\ n\} = A$
<proof>

lemma *add-0-left-vecset*:
assumes $(A :: 'a :: \text{monoid-add vec set}) \subseteq \text{carrier-vec } n$
shows $\{0_v\ n\} + A = A$
<proof>

lemma *assoc-add-vecset*:
assumes $(A :: 'a :: \text{semigroup-add vec set}) \subseteq \text{carrier-vec } n$
 and $B \subseteq \text{carrier-vec } n$
 and $C \subseteq \text{carrier-vec } n$
shows $A + (B + C) = (A + B) + C$
<proof>

lemma *sum-carrier-vec[intro]*: $A \subseteq \text{carrier-vec } n \implies B \subseteq \text{carrier-vec } n \implies A + B \subseteq \text{carrier-vec } n$
<proof>

lemma *comm-add-vecset*:
assumes $(A :: 'a :: \text{ab-semigroup-add vec set}) \subseteq \text{carrier-vec } n$
 and $B \subseteq \text{carrier-vec } n$

shows $A + B = B + A$
 ⟨*proof*⟩

end

6 Integral and Bounded Matrices and Vectors

We define notions of integral vectors and matrices and bounded vectors and matrices and prove some preservation lemmas. Moreover, we prove a bound on determinants.

theory *Integral-Bounded-Vectors*

imports

Missing-VS-Connect

Sum-Vec-Set

LLL-Basis-Reduction.Gram-Schmidt-2

begin

lemma *sq-norm-unit-vec[simp]*: **assumes** $i: i < n$

shows $\|unit-vec\ n\ i\|^2 = (1 :: 'a :: \{comm-ring-1, conjugatable-ring\})$
 ⟨*proof*⟩

definition *Ints-vec* (\mathbf{Z}_v) **where**

$\mathbf{Z}_v = \{x. \forall i < dim-vec\ x. x\ \$\ i \in \mathbf{Z}\}$

definition *indexed-Ints-vec* **where**

indexed-Ints-vec $I = \{x. \forall i < dim-vec\ x. i \in I \longrightarrow x\ \$\ i \in \mathbf{Z}\}$

lemma *indexed-Ints-vec-UNIV*: $\mathbf{Z}_v = indexed-Ints-vec\ UNIV$

⟨*proof*⟩

lemma *indexed-Ints-vec-subset*: $\mathbf{Z}_v \subseteq indexed-Ints-vec\ I$

⟨*proof*⟩

lemma *Ints-vec-vec-set*: $v \in \mathbf{Z}_v = (vec-set\ v \subseteq \mathbf{Z})$

⟨*proof*⟩

definition *Ints-mat* (\mathbf{Z}_m) **where**

$\mathbf{Z}_m = \{A. \forall i < dim-row\ A. \forall j < dim-col\ A. A\ \$\$ (i,j) \in \mathbf{Z}\}$

lemma *Ints-mat-elements-mat*: $A \in \mathbf{Z}_m = (elements-mat\ A \subseteq \mathbf{Z})$

⟨*proof*⟩

lemma *minus-in-Ints-vec-iff[simp]*: $(-x) \in \mathbf{Z}_v \longleftrightarrow (x :: 'a :: ring-1\ vec) \in \mathbf{Z}_v$

⟨*proof*⟩

lemma *minus-in-Ints-mat-iff[simp]*: $(-A) \in \mathbf{Z}_m \longleftrightarrow (A :: 'a :: ring-1\ mat) \in \mathbf{Z}_m$

$\langle proof \rangle$

lemma *Ints-vec-rows-Ints-mat*[simp]: $set (rows A) \subseteq \mathbb{Z}_v \longleftrightarrow A \in \mathbb{Z}_m$
 $\langle proof \rangle$

lemma *unit-vec-integral*[simp,intro]: $unit-vec n i \in \mathbb{Z}_v$
 $\langle proof \rangle$

lemma *diff-indexed-Ints-vec*:
 $x \in carrier-vec n \implies y \in carrier-vec n \implies x \in indexed-Ints-vec I \implies y \in indexed-Ints-vec I$
 $\implies x - y \in indexed-Ints-vec I$
 $\langle proof \rangle$

lemma *smult-indexed-Ints-vec*: $x \in \mathbb{Z} \implies v \in indexed-Ints-vec I \implies x \cdot_v v \in indexed-Ints-vec I$
 $\langle proof \rangle$

lemma *add-indexed-Ints-vec*:
 $x \in carrier-vec n \implies y \in carrier-vec n \implies x \in indexed-Ints-vec I \implies y \in indexed-Ints-vec I$
 $\implies x + y \in indexed-Ints-vec I$
 $\langle proof \rangle$

lemma (in *vec-space*) *lincomb-indexed-Ints-vec*: **assumes** $cI: \bigwedge x. x \in C \implies c x \in \mathbb{Z}$
and $C: C \subseteq carrier-vec n$
and $CI: C \subseteq indexed-Ints-vec I$
shows $lincomb c C \in indexed-Ints-vec I$
 $\langle proof \rangle$

definition *Bounded-vec* ($b :: 'a :: linordered-idom$) = $\{x . \forall i < dim-vec x . abs (x \$ i) \leq b\}$

lemma *Bounded-vec-vec-set*: $v \in Bounded-vec b \longleftrightarrow (\forall x \in vec-set v . abs x \leq b)$
 $\langle proof \rangle$

definition *Bounded-mat* ($b :: 'a :: linordered-idom$) =
 $\{A . (\forall i < dim-row A . \forall j < dim-col A . abs (A \$\$ (i,j)) \leq b)\}$

lemma *Bounded-mat-elements-mat*: $A \in Bounded-mat b \longleftrightarrow (\forall x \in elements-mat A . abs x \leq b)$
 $\langle proof \rangle$

lemma *Bounded-vec-rows-Bounded-mat*[simp]: $set (rows A) \subseteq Bounded-vec B \longleftrightarrow A \in Bounded-mat B$
 $\langle proof \rangle$

lemma *unit-vec-Bounded-vec*[simp,intro]: $unit-vec n i \in Bounded-vec (max 1 Bnd)$

<proof>

lemma *Bounded-matD*: **assumes** $A \in \text{Bounded-mat } b$
 $A \in \text{carrier-mat } nr \ nc$
shows $i < nr \implies j < nc \implies \text{abs } (A \ \$\$ (i,j)) \leq b$
<proof>

lemma *Bounded-vec-mono*: $b \leq B \implies \text{Bounded-vec } b \subseteq \text{Bounded-vec } B$
<proof>

lemma *Bounded-mat-mono*: $b \leq B \implies \text{Bounded-mat } b \subseteq \text{Bounded-mat } B$
<proof>

lemma *finite-Bounded-vec-Max*:
assumes $A \subseteq \text{carrier-vec } n$
and *fin*: *finite* A
shows $A \subseteq \text{Bounded-vec } (\text{Max } \{ \text{abs } (a \ \$ \ i) \mid a \ i. \ a \in A \wedge i < n \})$
<proof>

definition *det-bound* :: $\text{nat} \Rightarrow 'a :: \text{linordered-idom} \Rightarrow 'a$ **where**
 $\text{det-bound } n \ x = \text{fact } n * (x \hat{\ } n)$

lemma *det-bound*: **assumes** $A \in \text{carrier-mat } n \ n$
and $x: A \in \text{Bounded-mat } x$
shows $\text{abs } (\text{det } A) \leq \text{det-bound } n \ x$
<proof>

lemma *minus-in-Bounded-vec[simp]*:
 $(-x) \in \text{Bounded-vec } b \longleftrightarrow x \in \text{Bounded-vec } b$
<proof>

lemma *sum-in-Bounded-vecI[intro]*: **assumes**
 $xB: x \in \text{Bounded-vec } B1$ **and**
 $yB: y \in \text{Bounded-vec } B2$ **and**
 $x: x \in \text{carrier-vec } n$ **and**
 $y: y \in \text{carrier-vec } n$
shows $x + y \in \text{Bounded-vec } (B1 + B2)$
<proof>

lemma (**in** *gram-schmidt*) *lincomb-card-bound*: **assumes** $XBnd: X \subseteq \text{Bounded-vec } Bnd$
Bnd
and $X: X \subseteq \text{carrier-vec } n$
and $Bnd: Bnd \geq 0$
and $c: \bigwedge x. x \in X \implies \text{abs } (c \ x) \leq 1$
and $\text{card}: \text{card } X \leq k$
shows $\text{lincomb } c \ X \in \text{Bounded-vec } (\text{of-nat } k * Bnd)$
<proof>

lemma *bounded-vecset-sum*:

```

assumes  $A_{\text{carr}}$ :  $A \subseteq \text{carrier-vec } n$ 
and  $B_{\text{carr}}$ :  $B \subseteq \text{carrier-vec } n$ 
and  $\text{sum}$ :  $C = A + B$ 
and  $C_{\text{bnd}}$ :  $\exists \text{ bnd}C. C \subseteq \text{Bounded-vec bnd}C$ 
shows  $A \neq \{\}$   $\implies (\exists \text{ bnd}B. B \subseteq \text{Bounded-vec bnd}B)$ 
and  $B \neq \{\}$   $\implies (\exists \text{ bnd}A. A \subseteq \text{Bounded-vec bnd}A)$ 
<proof>

end

```

7 Cones

We define the notions like cone, polyhedral cone, etc. and prove some basic facts about them.

```

theory Cone
imports
  Basis-Extension
  Missing-VS-Connect
  Integral-Bounded-Vectors
begin

context gram-schmidt
begin

definition nonneg-lincomb  $c \ Vs \ b = (\text{lincomb } c \ Vs = b \wedge c \ ' \ Vs \subseteq \{x. x \geq 0\})$ 
definition nonneg-lincomb-list  $c \ Vs \ b = (\text{lincomb-list } c \ Vs = b \wedge (\forall i < \text{length } Vs. c \ i \geq 0))$ 

definition finite-cone  $:: 'a \ \text{vec set} \Rightarrow 'a \ \text{vec set}$  where
  finite-cone  $Vs = (\{ b. \exists c. \text{nonneg-lincomb } c \ (\text{if finite } Vs \ \text{then } Vs \ \text{else } \{\}) \ b})$ 

definition cone  $:: 'a \ \text{vec set} \Rightarrow 'a \ \text{vec set}$  where
  cone  $Vs = (\{ x. \exists Ws. \text{finite } Ws \wedge Ws \subseteq Vs \wedge x \in \text{finite-cone } Ws})$ 

definition cone-list  $:: 'a \ \text{vec list} \Rightarrow 'a \ \text{vec set}$  where
  cone-list  $Vs = \{b. \exists c. \text{nonneg-lincomb-list } c \ Vs \ b\}$ 

lemma finite-cone-iff-cone-list: assumes  $Vs: Vs \subseteq \text{carrier-vec } n$ 
and  $\text{id}: Vs = \text{set } Vsl$ 
shows finite-cone  $Vs = \text{cone-list } Vsl$ 
<proof>

lemma cone-alt-def: assumes  $Vs: Vs \subseteq \text{carrier-vec } n$ 
shows cone  $Vs = (\{ x. \exists Ws. \text{set } Ws \subseteq Vs \wedge x \in \text{cone-list } Ws})$ 
<proof>

lemma cone-mono:  $Vs \subseteq Ws \implies \text{cone } Vs \subseteq \text{cone } Ws$ 
<proof>

```


lemma *finite-cone-mono*: **assumes** *fin*: *finite* *Ws*
and *Ws*: $Ws \subseteq \text{carrier-vec } n$
and *sub*: $Vs \subseteq Ws$
shows *finite-cone* $Vs \subseteq \text{finite-cone } Ws$
<proof>

lemma *finite-cone-carrier*: $A \subseteq \text{carrier-vec } n \implies \text{finite-cone } A \subseteq \text{carrier-vec } n$
<proof>

lemma *cone-carrier*: $A \subseteq \text{carrier-vec } n \implies \text{cone } A \subseteq \text{carrier-vec } n$
<proof>

lemma *cone-iff-finite-cone*: **assumes** *A*: $A \subseteq \text{carrier-vec } n$
and *fin*: *finite* *A*
shows $\text{cone } A = \text{finite-cone } A$
<proof>

lemma *set-in-finite-cone*:
assumes *Vs*: $Vs \subseteq \text{carrier-vec } n$
and *fin*: *finite* *Vs*
shows $Vs \subseteq \text{finite-cone } Vs$
<proof>

lemma *set-in-cone*:
assumes *Vs*: $Vs \subseteq \text{carrier-vec } n$
shows $Vs \subseteq \text{cone } Vs$
<proof>

lemma *zero-in-finite-cone*:
assumes *Vs*: $Vs \subseteq \text{carrier-vec } n$
shows $0_v \in \text{finite-cone } Vs$
<proof>

lemma *lincomb-in-finite-cone*:
assumes *x* = *lincomb* *l* *W*
and *finite* *W*
and $\forall i \in W . l\ i \geq 0$
and $W \subseteq \text{carrier-vec } n$
shows $x \in \text{finite-cone } W$
<proof>

lemma *lincomb-in-cone*:
assumes *x* = *lincomb* *l* *W*
and *finite* *W*
and $\forall i \in W . l\ i \geq 0$
and $W \subseteq \text{carrier-vec } n$
shows $x \in \text{cone } W$
<proof>

lemma *zero-in-cone*: $0_v \ n \in \text{cone } Vs$
(*proof*)

lemma *cone-smult*:
 assumes $a: a \geq 0$
 and $Vs: Vs \subseteq \text{carrier-vec } n$
 and $x: x \in \text{cone } Vs$
 shows $a \cdot_v x \in \text{cone } Vs$
(*proof*)

lemma *finite-cone-empty[simp]*: $\text{finite-cone } \{\} = \{0_v \ n\}$
(*proof*)

lemma *cone-empty[simp]*: $\text{cone } \{\} = \{0_v \ n\}$
(*proof*)

lemma *cone-elem-sum*:
 assumes $Vs: Vs \subseteq \text{carrier-vec } n$
 and $x: x \in \text{cone } Vs$
 and $y: y \in \text{cone } Vs$
 shows $x + y \in \text{cone } Vs$
(*proof*)

lemma *cone-cone*:
 assumes $Vs: Vs \subseteq \text{carrier-vec } n$
 shows $\text{cone } (\text{cone } Vs) = \text{cone } Vs$
(*proof*)

lemma *cone-smult-basis*:
 assumes $Vs: Vs \subseteq \text{carrier-vec } n$
 and $l: l \cdot Vs \subseteq \{x. x > 0\}$
 shows $\text{cone } \{l \cdot v \cdot_v v \mid v \cdot v \in Vs\} = \text{cone } Vs$
(*proof*)

lemma *cone-add-cone*:
 assumes $C: C \subseteq \text{carrier-vec } n$
 shows $\text{cone } C + \text{cone } C = \text{cone } C$
(*proof*)

lemma *orthogonal-cone*:
 assumes $X: X \subseteq \text{carrier-vec } n$
 and $W: W \subseteq \text{carrier-vec } n$
 and $\text{fin}X: \text{finite } X$
 and $\text{span}LW: \text{span } (\text{set } Ls \cup W) = \text{carrier-vec } n$
 and $\text{ortho}: \bigwedge w \ x. w \in W \implies x \in \text{set } Ls \implies w \cdot x = 0$
 and $WWs: W = \text{set } Ws$
 and $\text{span}L: \text{span } (\text{set } Ls) = \text{span } X$

and LX : $set\ Ls \subseteq X$
and $lin\text{-}Ls\text{-}Bs$: $lin\text{-}indpt\text{-}list\ (Ls\ @\ Bs)$
and $len\text{-}Ls\text{-}Bs$: $length\ (Ls\ @\ Bs) = n$
shows $cone\ (X \cup\ set\ Bs) \cap\ \{x \in\ carrier\text{-}vec\ n.\ \forall\ w \in\ W.\ w \cdot\ x = 0\} = cone\ X$
 $\wedge\ x.\ \forall\ w \in\ W.\ w \cdot\ x = 0 \implies Z \subseteq X \implies B \subseteq set\ Bs \implies x = lincomb\ c\ (Z \cup$
 $B)$
 $\implies x = lincomb\ c\ (Z - B)$
 $\langle proof \rangle$

definition $polyhedral\text{-}cone\ (A :: 'a\ mat) = \{x.\ x \in\ carrier\text{-}vec\ n \wedge A *_{\nu}\ x \leq 0_{\nu}$
 $(dim\text{-}row\ A)\}$

lemma $polyhedral\text{-}cone\text{-}carrier$: **assumes** $A \in\ carrier\text{-}mat\ nr\ n$
shows $polyhedral\text{-}cone\ A \subseteq carrier\text{-}vec\ n$
 $\langle proof \rangle$

lemma $cone\text{-}in\text{-}polyhedral\text{-}cone$:
assumes CA : $C \subseteq polyhedral\text{-}cone\ A$
and A : $A \in\ carrier\text{-}mat\ nr\ n$
shows $cone\ C \subseteq polyhedral\text{-}cone\ A$
 $\langle proof \rangle$

lemma $bounded\text{-}cone\text{-}is\text{-}zero$:
assumes $Carr$: $C \subseteq carrier\text{-}vec\ n$ **and** bnd : $cone\ C \subseteq Bounded\text{-}vec\ bnd$
shows $cone\ C = \{0_{\nu}\ n\}$
 $\langle proof \rangle$

lemma $cone\text{-}of\text{-}cols$: **fixes** $A :: 'a\ mat$ **and** $b :: 'a\ vec$
assumes A : $A \in\ carrier\text{-}mat\ n\ nr$ **and** b : $b \in\ carrier\text{-}vec\ n$
shows $b \in\ cone\ (set\ (cols\ A)) \longleftrightarrow (\exists\ x.\ x \geq 0_{\nu}\ nr \wedge A *_{\nu}\ x = b)$
 $\langle proof \rangle$

end
end

8 Convex Hulls

We define the notion of convex hull of a set or list of vectors and derive basic properties thereof.

theory $Convex\text{-}Hull$
imports $Cone$
begin

context $gram\text{-}schmidt$
begin

definition $convex\text{-}lincomb\ c\ Vs\ b = (nonneg\text{-}lincomb\ c\ Vs\ b \wedge sum\ c\ Vs = 1)$

definition *convex-lincomb-list* c Vs $b = (\text{nonneg-lincomb-list } c \text{ } Vs \text{ } b \wedge \text{sum } c \{0..<\text{length } Vs\} = 1)$

definition *convex-hull* $Vs = \{x. \exists \text{ } Ws \text{ } c. \text{finite } Ws \wedge Ws \subseteq Vs \wedge \text{convex-lincomb } c \text{ } Ws \text{ } x\}$

lemma *convex-hull-carrier[intro]*: $Vs \subseteq \text{carrier-vec } n \implies \text{convex-hull } Vs \subseteq \text{carrier-vec } n$
 ⟨proof⟩

lemma *convex-hull-mono*: $Vs \subseteq Ws \implies \text{convex-hull } Vs \subseteq \text{convex-hull } Ws$
 ⟨proof⟩

lemma *convex-lincomb-empty[simp]*: $\neg (\text{convex-lincomb } c \text{ } \{\} \text{ } x)$
 ⟨proof⟩

lemma *set-in-convex-hull*:
 assumes $A \subseteq \text{carrier-vec } n$
 shows $A \subseteq \text{convex-hull } A$
 ⟨proof⟩

lemma *convex-hull-empty[simp]*:
 $\text{convex-hull } \{\} = \{\}$
 $A \subseteq \text{carrier-vec } n \implies \text{convex-hull } A = \{\} \longleftrightarrow A = \{\}$
 ⟨proof⟩

lemma *convex-hull-bound*: **assumes** $XBnd$: $X \subseteq \text{Bounded-vec } Bnd$
and $X: X \subseteq \text{carrier-vec } n$
shows $\text{convex-hull } X \subseteq \text{Bounded-vec } Bnd$
 ⟨proof⟩

definition *convex-hull-list* $Vs = \{x. \exists \text{ } c. \text{convex-lincomb-list } c \text{ } Vs \text{ } x\}$

lemma *lincomb-list-elem*:
 $\text{set } Vs \subseteq \text{carrier-vec } n \implies$
 $\text{lincomb-list } (\lambda j. \text{if } i=j \text{ then } 1 \text{ else } 0) \text{ } Vs = (\text{if } i < \text{length } Vs \text{ then } Vs ! i \text{ else } 0_v$
 $n)$
 ⟨proof⟩

lemma *set-in-convex-hull-list*: **fixes** $Vs :: 'a \text{ vec list}$
assumes $\text{set } Vs \subseteq \text{carrier-vec } n$
shows $\text{set } Vs \subseteq \text{convex-hull-list } Vs$
 ⟨proof⟩

lemma *convex-hull-list-combination*:
assumes $Vs: \text{set } Vs \subseteq \text{carrier-vec } n$
and $x: x \in \text{convex-hull-list } Vs$
and $y: y \in \text{convex-hull-list } Vs$
and $l0: 0 \leq l$ **and** $l1: l \leq 1$

shows $l \cdot_v x + (1 - l) \cdot_v y \in \text{convex-hull-list } Vs$
 ⟨proof⟩

lemma *convex-hull-list-mono*:

assumes $set\ Vs \subseteq \text{carrier-vec } n$

shows $set\ Vs \subseteq set\ Ws \implies \text{convex-hull-list } Vs \subseteq \text{convex-hull-list } Ws$
 ⟨proof⟩

lemma *convex-hull-list-eq-set*:

set $Vs \subseteq \text{carrier-vec } n \implies set\ Vs = set\ Ws \implies \text{convex-hull-list } Vs = \text{convex-hull-list } Ws$
 ⟨proof⟩

lemma *find-indices-empty*: $(\text{find-indices } x\ Vs = []) = (x \notin set\ Vs)$
 ⟨proof⟩

lemma *distinct-list-find-indices*:

shows $\llbracket i < \text{length } Vs; Vs\ !\ i = x; \text{distinct } Vs \rrbracket \implies \text{find-indices } x\ Vs = [i]$
 ⟨proof⟩

lemma *finite-convex-hull-iff-convex-hull-list*: **assumes** $Vs: Vs \subseteq \text{carrier-vec } n$

and $id': Vs = set\ Vsl'$

shows $\text{convex-hull } Vs = \text{convex-hull-list } Vsl'$
 ⟨proof⟩

definition $\text{convex } S = (\text{convex-hull } S = S)$

lemma *convex-convex-hull*: $\text{convex } S \implies \text{convex-hull } S = S$
 ⟨proof⟩

lemma *convex-hull-convex-hull-listD*: **assumes** $A: A \subseteq \text{carrier-vec } n$

and $x: x \in \text{convex-hull } A$

shows $\exists as. set\ as \subseteq A \wedge x \in \text{convex-hull-list } as$
 ⟨proof⟩

lemma *convex-hull-convex-sum*: **assumes** $A: A \subseteq \text{carrier-vec } n$

and $x: x \in \text{convex-hull } A$

and $y: y \in \text{convex-hull } A$

and $a: 0 \leq a \wedge a \leq 1$

shows $a \cdot_v x + (1 - a) \cdot_v y \in \text{convex-hull } A$
 ⟨proof⟩

lemma *convexI*: **assumes** $S: S \subseteq \text{carrier-vec } n$

and $step: \bigwedge a\ x\ y. x \in S \implies y \in S \implies 0 \leq a \implies a \leq 1 \implies a \cdot_v x + (1 - a) \cdot_v y \in S$

shows $\text{convex } S$
 ⟨proof⟩

lemma *convex-hulls-are-convex*: **assumes** $A \subseteq \text{carrier-vec } n$

```

shows convex (convex-hull A)
  ⟨proof⟩

lemma convex-hull-sum: assumes A: A ⊆ carrier-vec n and B: B ⊆ carrier-vec
n
  shows convex-hull (A + B) = convex-hull A + convex-hull B
  ⟨proof⟩

lemma convex-hull-in-cone:
  convex-hull C ⊆ cone C
  ⟨proof⟩

lemma convex-cone:
  assumes C: C ⊆ carrier-vec n
  shows convex (cone C)
  ⟨proof⟩

end
end

```

9 Normal Vectors

We provide a function for the normal vector of a half-space (given as $n-1$ linearly independent vectors). We further provide a function that returns a list of normal vectors that span the orthogonal complement of some subspace of R^n . Bounds for all normal vectors are provided.

```

theory Normal-Vector
  imports
    Integral-Bounded-Vectors
    Basis-Extension
  begin

  context gram-schmidt
  begin

  lemma ortho-sum-in-span:
    assumes W: W ⊆ carrier-vec n
      and X: X ⊆ carrier-vec n
      and ortho:  $\bigwedge w x. w \in W \implies x \in X \implies x \cdot w = 0$ 
      and inspan:  $\text{lincomb } l1 X + \text{lincomb } l2 W \in \text{span } X$ 
    shows  $\text{lincomb } l2 W = 0_v n$ 
    ⟨proof⟩

  lemma ortho-lin-indpt: assumes W: W ⊆ carrier-vec n
    and X: X ⊆ carrier-vec n
    and ortho:  $\bigwedge w x. w \in W \implies x \in X \implies x \cdot w = 0$ 
    and linW: lin-indpt W

```

and *linX*: *lin-indpt X*
shows *lin-indpt (W ∪ X)*
 ⟨*proof*⟩

definition *normal-vector* :: 'a *vec set* ⇒ 'a *vec* **where**
normal-vector W = (let *ws* = (SOME *ws. set ws = W ∧ distinct ws*);
m = *length ws*;
B = (λ *j. mat m m (λ(i, j'). ws ! i \$ (if j' < j then j' else Suc j'))*))
 in *vec n (λ j. (-1)^(m+j) * det (B j))*)

lemma *normal-vector: assumes fin: finite W*
and *card: Suc (card W) = n*
and *lin: lin-indpt W*
and *W: W ⊆ carrier-vec n*
shows *normal-vector W ∈ carrier-vec n*
normal-vector W ≠ 0_v n
w ∈ W ⇒ w · normal-vector W = 0
w ∈ W ⇒ normal-vector W · w = 0
lin-indpt (insert (normal-vector W) W)
normal-vector W ∉ W
W ⊆ Z_v ∩ Bounded-vec Bnd ⇒ normal-vector W ∈ Z_v ∩ Bounded-vec (det-bound (n-1) Bnd)
 ⟨*proof*⟩

lemma *normal-vector-span:*
assumes *card: Suc (card D) = n*
and *D: D ⊆ carrier-vec n* **and** *fin: finite D* **and** *lin: lin-indpt D*
shows *span D = { x. x ∈ carrier-vec n ∧ x · normal-vector D = 0 }*
 ⟨*proof*⟩

definition *normal-vectors* :: 'a *vec list* ⇒ 'a *vec list* **where**
normal-vectors ws = (let *us* = *basis-extension ws*
 in *map (λ i. normal-vector (set (ws @ us) - {us ! i})) [0..<length us]*)

lemma *normal-vectors:*
assumes *lin: lin-indpt-list ws*
shows *set (normal-vectors ws) ⊆ carrier-vec n*
w ∈ set ws ⇒ nv ∈ set (normal-vectors ws) ⇒ nv · w = 0
w ∈ set ws ⇒ nv ∈ set (normal-vectors ws) ⇒ w · nv = 0
lin-indpt-list (ws @ normal-vectors ws)
length ws + length (normal-vectors ws) = n
set ws ∩ set (normal-vectors ws) = {}
set ws ⊆ Z_v ∩ Bounded-vec Bnd ⇒
set (normal-vectors ws) ⊆ Z_v ∩ Bounded-vec (det-bound (n-1) (max 1 Bnd))
 ⟨*proof*⟩

definition *pos-norm-vec* :: 'a *vec set* ⇒ 'a *vec* ⇒ 'a *vec* **where**
pos-norm-vec D x = (let *c'* = *normal-vector D*;
c = (if *c' · x > 0* then *c'* else *-c'*) in *c*)

lemma *pos-norm-vec*:
assumes $D: D \subseteq \text{carrier-vec } n$ **and** $\text{fin}: \text{finite } D$ **and** $\text{lin}: \text{lin-indpt } D$
and $\text{card}: \text{Suc } (\text{card } D) = n$
and $c\text{-def}: c = \text{pos-norm-vec } D \ x$
shows $c \in \text{carrier-vec } n \ \text{span } D = \{x. x \in \text{carrier-vec } n \wedge x \cdot c = 0\}$
 $x \notin \text{span } D \implies x \in \text{carrier-vec } n \implies c \cdot x > 0$
 $c \in \{\text{normal-vector } D, -\text{normal-vector } D\}$
 $\langle \text{proof} \rangle$

end

end

10 Dimension of Spans

We define the notion of dimension of a span of vectors and prove some natural results about them. The definition is made as a function, so that no interpretation of locales like subspace is required.

theory *Dim-Span*
imports *Missing-VS-Connect*
begin

context *vec-space*

begin

definition $\text{dim-span } W = \text{Max } (\text{card } \{V. V \subseteq \text{carrier-vec } n \wedge V \subseteq \text{span } W \wedge \text{lin-indpt } V\})$

lemma **fixes** $V \ W :: 'a \ \text{vec set}$

shows

card-le-dim-span:

$V \subseteq \text{carrier-vec } n \implies V \subseteq \text{span } W \implies \text{lin-indpt } V \implies \text{card } V \leq \text{dim-span } W$

and

card-eq-dim-span-imp-same-span:

$W \subseteq \text{carrier-vec } n \implies V \subseteq \text{span } W \implies \text{lin-indpt } V \implies \text{card } V = \text{dim-span } W$

$W \implies \text{span } V = \text{span } W$ **and**

same-span-imp-card-eq-dim-span:

$V \subseteq \text{carrier-vec } n \implies W \subseteq \text{carrier-vec } n \implies \text{span } V = \text{span } W \implies \text{lin-indpt } V \implies \text{card } V = \text{dim-span } W$ **and**

dim-span-cong:

$\text{span } V = \text{span } W \implies \text{dim-span } V = \text{dim-span } W$ **and**

ex-basis-span:

$V \subseteq \text{carrier-vec } n \implies \exists W. W \subseteq \text{carrier-vec } n \wedge \text{lin-indpt } W \wedge \text{span } V = \text{span } W \wedge \text{dim-span } V = \text{card } W$

$\langle \text{proof} \rangle$

lemma *dim-span-le-n*: **assumes** $W: W \subseteq \text{carrier-vec } n$ **shows** $\text{dim-span } W \leq n$
 $\langle \text{proof} \rangle$

lemma *dim-span-insert*: **assumes** $W: W \subseteq \text{carrier-vec } n$
and $v: v \in \text{carrier-vec } n$ **and** $vs: v \notin \text{span } W$
shows $\text{dim-span } (\text{insert } v \ W) = \text{Suc } (\text{dim-span } W)$
 $\langle \text{proof} \rangle$
end
end

11 The Fundamental Theorem of Linear Inequalities

The theorem states that for a given set of vectors A and vector b , either b is in the cone of a linear independent subset of A , or there is a hyperplane that contains $\text{span}(A,b)-1$ linearly independent vectors of A that separates A from b . We prove this theorem and derive some consequences, e.g., Caratheodory's theorem that b is the cone of A iff b is in the cone of a linear independent subset of A .

theory *Fundamental-Theorem-Linear-Inequalities*

imports

Cone

Normal-Vector

Dim-Span

begin

context *gram-schmidt*

begin

The mentions equivances A-D are:

- A: b is in the cone of vectors A ,
- B: b is in the cone of a subset of linear independent of vectors A ,
- C: there is no separating hyperplane of b and the vectors A , which contains dim many linear independent vectors of A
- D: there is no separating hyperplane of b and the vectors A

lemma *fundamental-theorem-of-linear-inequalities-A-imp-D*:

assumes $A: A \subseteq \text{carrier-vec } n$

and $fin: \text{finite } A$

and $b: b \in \text{cone } A$

shows $\nexists c. c \in \text{carrier-vec } n \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0$

$\langle \text{proof} \rangle$

The difficult direction is that C implies B. To this end we follow the proof that at least one of B and the negation of C is satisfied.

context

fixes $a :: \text{nat} \Rightarrow 'a \text{ vec}$
and $b :: 'a \text{ vec}$
and $m :: \text{nat}$
assumes $a : a \text{ ' } \{0 \dots m\} \subseteq \text{carrier-vec } n$
and $\text{inj-}a : \text{inj-on } a \{0 \dots m\}$
and $b : b \in \text{carrier-vec } n$
and $\text{full-span} : \text{span } (a \text{ ' } \{0 \dots m\}) = \text{carrier-vec } n$

begin

private definition $\text{goal} = ((\exists I. I \subseteq \{0 \dots m\} \wedge \text{card } (a \text{ ' } I) = n \wedge \text{lin-indpt}$
 $(a \text{ ' } I) \wedge b \in \text{finite-cone } (a \text{ ' } I))$
 $\vee (\exists c I. I \subseteq \{0 \dots m\} \wedge c \in \{\text{normal-vector } (a \text{ ' } I), - \text{normal-vector } (a \text{ ' } I)\}$
 $\wedge \text{Suc } (\text{card } (a \text{ ' } I)) = n$
 $\wedge \text{lin-indpt } (a \text{ ' } I) \wedge (\forall i < m. c \cdot a \ i \geq 0) \wedge c \cdot b < 0))$

private lemma $\text{card-}a\text{-}I[\text{simp}] : I \subseteq \{0 \dots m\} \implies \text{card } (a \text{ ' } I) = \text{card } I$
 $\langle \text{proof} \rangle$ **lemma** $\text{in-}a\text{-}I[\text{simp}] : I \subseteq \{0 \dots m\} \implies i < m \implies (a \ i \in a \text{ ' } I) = (i$
 $\in I)$
 $\langle \text{proof} \rangle$ **definition** $\text{valid-}I = \{ I. \text{card } I = n \wedge \text{lin-indpt } (a \text{ ' } I) \wedge I \subseteq \{0 \dots$
 $m\} \}$

private definition $\text{cond where } \text{cond } I \ I' \ l \ c \ h \ k \equiv$
 $b = \text{lincomb } l \ (a \text{ ' } I) \wedge$
 $h \in I \wedge l \ (a \ h) < 0 \wedge (\forall h'. h' \in I \longrightarrow h' < h \longrightarrow l \ (a \ h') \geq 0) \wedge$
 $c \in \text{carrier-vec } n \wedge \text{span } (a \text{ ' } (I - \{h\})) = \{ x. x \in \text{carrier-vec } n \wedge c \cdot x = 0 \}$
 $\wedge c \cdot b < 0 \wedge$
 $k < m \wedge c \cdot a \ k < 0 \wedge (\forall k'. k' < k \longrightarrow c \cdot a \ k' \geq 0) \wedge$
 $I' = \text{insert } k \ (I - \{h\})$

private definition $\text{step-rel} = \text{Restr } \{ (I'', I). \exists l \ c \ h \ k. \text{cond } I \ I'' \ l \ c \ h \ k \} \ \text{valid-}I$

private lemma $\text{finite-step-rel} : \text{finite step-rel}$
 $\langle \text{proof} \rangle$ **lemma** $\text{acyclic-imp-goal} : \text{acyclic step-rel} \implies \text{goal}$
 $\langle \text{proof} \rangle$ **lemma** $\text{acyclic-step-rel} : \text{acyclic step-rel}$
 $\langle \text{proof} \rangle$

lemma $\text{fundamental-theorem-neg-C-or-B-in-context} :$

assumes $W : W = a \text{ ' } \{0 \dots m\}$
shows $(\exists U. U \subseteq W \wedge \text{card } U = n \wedge \text{lin-indpt } U \wedge b \in \text{finite-cone } U) \vee$
 $(\exists c U. U \subseteq W \wedge$
 $c \in \{\text{normal-vector } U, - \text{normal-vector } U\} \wedge$
 $\text{Suc } (\text{card } U) = n \wedge \text{lin-indpt } U \wedge (\forall w \in W. 0 \leq c \cdot w) \wedge c \cdot b < 0)$
 $\langle \text{proof} \rangle$

end

lemma $\text{fundamental-theorem-of-linear-inequalities-C-imp-B-full-dim} :$

assumes $A : A \subseteq \text{carrier-vec } n$

and *fin*: *finite A*
and *span*: $\text{span } A = \text{carrier-vec } n$
and *b*: $b \in \text{carrier-vec } n$
and *C*: $\exists c \in B. B \subseteq A \wedge c \in \{\text{normal-vector } B, - \text{normal-vector } B\} \wedge \text{Suc}(\text{card } B) = n$
 $\wedge \text{lin-indpt } B \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0$
shows $\exists B \subseteq A. \text{lin-indpt } B \wedge \text{card } B = n \wedge b \in \text{finite-cone } B$
<proof>

lemma *fundamental-theorem-of-linear-inequalities-full-dim*: **fixes** *A* :: 'a vec set
defines *HyperN* $\equiv \{b. b \in \text{carrier-vec } n \wedge (\exists B \subseteq A. B \subseteq A \wedge c \in \{\text{normal-vector } B, - \text{normal-vector } B\} \wedge \text{Suc}(\text{card } B) = n \wedge \text{lin-indpt } B \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0)\}$
defines *HyperA* $\equiv \{b. b \in \text{carrier-vec } n \wedge (\exists c. c \in \text{carrier-vec } n \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0)\}$
defines *lin-indpt-cone* $\equiv \bigcup \{\text{finite-cone } B \mid B. B \subseteq A \wedge \text{card } B = n \wedge \text{lin-indpt } B\}$
assumes *A*: $A \subseteq \text{carrier-vec } n$
and *fin*: *finite A*
and *span*: $\text{span } A = \text{carrier-vec } n$
shows
 $\text{cone } A = \text{lin-indpt-cone}$
 $\text{cone } A = \text{HyperN}$
 $\text{cone } A = \text{HyperA}$
<proof>

lemma *fundamental-theorem-of-linear-inequalities-C-imp-B*:
assumes *A*: $A \subseteq \text{carrier-vec } n$
and *fin*: *finite A*
and *b*: $b \in \text{carrier-vec } n$
and *C*: $\exists c \in A'. c \in \text{carrier-vec } n$
 $\wedge A' \subseteq A \wedge \text{Suc}(\text{card } A') = \text{dim-span } (\text{insert } b \ A)$
 $\wedge (\forall a \in A'. c \cdot a = 0)$
 $\wedge \text{lin-indpt } A' \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0$
shows $\exists B \subseteq A. \text{lin-indpt } B \wedge \text{card } B = \text{dim-span } A \wedge b \in \text{finite-cone } B$
<proof>

lemma *fundamental-theorem-of-linear-inequalities*: **fixes** *A* :: 'a vec set
defines *HyperN* $\equiv \{b. b \in \text{carrier-vec } n \wedge (\exists c \in B. c \in \text{carrier-vec } n \wedge B \subseteq A \wedge \text{Suc}(\text{card } B) = \text{dim-span } (\text{insert } b \ A) \wedge \text{lin-indpt } B \wedge (\forall a \in B. c \cdot a = 0) \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0)\}$
defines *HyperA* $\equiv \{b. b \in \text{carrier-vec } n \wedge (\exists c. c \in \text{carrier-vec } n \wedge (\forall a_i \in A. c \cdot a_i \geq 0) \wedge c \cdot b < 0)\}$
defines *lin-indpt-cone* $\equiv \bigcup \{\text{finite-cone } B \mid B. B \subseteq A \wedge \text{card } B = \text{dim-span } A \wedge \text{lin-indpt } B\}$
assumes *A*: $A \subseteq \text{carrier-vec } n$
and *fin*: *finite A*

shows
cone $A = \text{lin-indpt-cone}$
cone $A = \text{Hyper}N$
cone $A = \text{Hyper}A$
 ⟨*proof*⟩

corollary *Caratheodory-theorem*: **assumes** $A: A \subseteq \text{carrier-vec } n$
shows $\text{cone } A = \bigcup \{ \text{finite-cone } B \mid B. B \subseteq A \wedge \text{lin-indpt } B \}$
 ⟨*proof*⟩
end
end

12 Farkas' Lemma

We prove two variants of Farkas' lemma. Note that type here is more general than in the versions of Farkas' Lemma which are in the AFP-entry Farkas-Lemma, which is restricted to rational matrices. However, there δ -rationals are supported, which are not present here.

theory *Farkas-Lemma*
imports *Fundamental-Theorem-Linear-Inequalities*
begin

context *gram-schmidt*
begin

lemma *Farkas-Lemma*: **fixes** $A :: 'a \text{ mat}$ **and** $b :: 'a \text{ vec}$
assumes $A: A \in \text{carrier-mat } n \text{ nr}$ **and** $b: b \in \text{carrier-vec } n$
shows $(\exists x. x \geq 0_v \text{ nr} \wedge A *_{\text{v}} x = b) \iff (\forall y. y \in \text{carrier-vec } n \longrightarrow A^T *_{\text{v}} y \geq 0_v \text{ nr} \longrightarrow y \cdot b \geq 0)$
 ⟨*proof*⟩

lemma *Farkas-Lemma'*:
fixes $A :: 'a \text{ mat}$ **and** $b :: 'a \text{ vec}$
assumes $A: A \in \text{carrier-mat } nr \text{ nc}$ **and** $b: b \in \text{carrier-vec } nr$
shows $(\exists x. x \in \text{carrier-vec } nc \wedge A *_{\text{v}} x \leq b)$
 $\iff (\forall y. y \geq 0_v \text{ nr} \wedge A^T *_{\text{v}} y = 0_v \text{ nc} \longrightarrow y \cdot b \geq 0)$
 ⟨*proof*⟩

end
end

13 The Theorem of Farkas, Minkowsky and Weyl

We prove the theorem of Farkas, Minkowsky and Weyl that a cone is finitely generated iff it is polyhedral. Moreover, we provide quantitative bounds.

theory *Farkas-Minkowsky-Weyl*
imports *Fundamental-Theorem-Linear-Inequalities*

begin

context *gram-schmidt*

begin

We first prove the one direction of the theorem for the case that the span of the vectors is the full n-dimensional space.

lemma *farkas-minkowsky-weyl-theorem-1-full-dim:*

assumes $X: X \subseteq \text{carrier-vec } n$

and $\text{fin}: \text{finite } X$

and $\text{span}: \text{span } X = \text{carrier-vec } n$

shows $\exists nr A. A \in \text{carrier-mat } nr n \wedge \text{cone } X = \text{polyhedral-cone } A$

$\wedge (X \subseteq \mathbb{Z}_v \cap \text{Bounded-vec } Bnd \longrightarrow A \in \mathbb{Z}_m \cap \text{Bounded-mat } (\text{det-bound } (n-1) Bnd))$

<proof>

We next generalize the theorem to the case where X does not span the full space. To this end, we extend X by unit-vectors until the full space is spanned, and then add the normal-vectors of these unit-vectors which are orthogonal to span X as additional constraints to the resulting matrix.

lemma *farkas-minkowsky-weyl-theorem-1:*

assumes $X: X \subseteq \text{carrier-vec } n$

and $\text{fin}X: \text{finite } X$

shows $\exists nr A. A \in \text{carrier-mat } nr n \wedge \text{cone } X = \text{polyhedral-cone } A \wedge$

$(X \subseteq \mathbb{Z}_v \cap \text{Bounded-vec } Bnd \longrightarrow A \in \mathbb{Z}_m \cap \text{Bounded-mat } (\text{det-bound } (n-1) (\text{max } 1 Bnd)))$

<proof>

Now for the other direction.

lemma *farkas-minkowsky-weyl-theorem-2:*

assumes $A: A \in \text{carrier-mat } nr n$

shows $\exists X. X \subseteq \text{carrier-vec } n \wedge \text{finite } X \wedge \text{polyhedral-cone } A = \text{cone } X$

$\wedge (A \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd \longrightarrow X \subseteq \mathbb{Z}_v \cap \text{Bounded-vec } (\text{det-bound } (n-1) (\text{max } 1 Bnd)))$

<proof>

lemma *farkas-minkowsky-weyl-theorem:*

$(\exists X. X \subseteq \text{carrier-vec } n \wedge \text{finite } X \wedge P = \text{cone } X)$

$\longleftrightarrow (\exists A nr. A \in \text{carrier-mat } nr n \wedge P = \text{polyhedral-cone } A)$

<proof>

end

end

14 The Decomposition Theorem

This theory contains a proof of the fact, that every polyhedron can be decomposed into a convex hull of a finite set of points + a finitely generated

cone, including bounds on the numbers that are required in the decomposition. We further prove the inverse direction of this theorem (without bounds) and as a corollary, we derive that a polyhedron is bounded iff it is the convex hull of finitely many points, i.e., a polytope.

theory *Decomposition-Theorem*

imports

Farkas-Minkowsky-Weyl

Convex-Hull

begin

context *gram-schmidt*

begin

definition *polytope* $P = (\exists V. V \subseteq \text{carrier-vec } n \wedge \text{finite } V \wedge P = \text{convex-hull } V)$

definition *polyhedron* $A \ b = \{x \in \text{carrier-vec } n. A *_{\vee} x \leq b\}$

lemma *polyhedra-are-convex:*

assumes $A: A \in \text{carrier-mat } nr \ n$

and $b: b \in \text{carrier-vec } nr$

and $P: P = \text{polyhedron } A \ b$

shows *convex* P

<proof>

end

locale *gram-schmidt-m = n: gram-schmidt n f-ty + m: gram-schmidt m f-ty*

for $n \ m :: \text{nat}$ **and** *f-ty*

begin

lemma *vec-first-lincomb-list:*

assumes $Xs: \text{set } Xs \subseteq \text{carrier-vec } n$

and $nm: m \leq n$

shows $\text{vec-first } (n.\text{lincomb-list } c \ Xs) \ m =$
 $m.\text{lincomb-list } c \ (\text{map } (\lambda v. \text{vec-first } v \ m) \ Xs)$

<proof>

lemma *convex-hull-next-dim:*

assumes $n = m + 1$

and $X: X \subseteq \text{carrier-vec } n$

and *finite* X

and $Xm1: \forall y \in X. y \ \$ \ m = 1$

and $y\text{-dim}: y \in \text{carrier-vec } n$

and $y: y \ \$ \ m = 1$

shows $(\text{vec-first } y \ m \in m.\text{convex-hull } \{\text{vec-first } y \ m \mid y. y \in X\}) = (y \in n.\text{cone } X)$

<proof>

lemma *cone-next-dim:*

assumes $n = m + 1$

and $X: X \subseteq \text{carrier-vec } n$

and *finite* X

and $Xm0: \forall y \in X. y \$ m = 0$

and $y\text{-dim}: y \in \text{carrier-vec } n$

and $y: y \$ m = 0$

shows $(\text{vec-first } y \ m \in m.\text{cone } \{\text{vec-first } y \ m \mid y. y \in X\}) = (y \in n.\text{cone } X)$

<proof>

end

context *gram-schmidt*

begin

lemma *decomposition-theorem-polyhedra-1:*

assumes $A: A \in \text{carrier-mat } nr \ n$

and $b: b \in \text{carrier-vec } nr$

and $P: P = \text{polyhedron } A \ b$

shows $\exists Q \ X. X \subseteq \text{carrier-vec } n \wedge \text{finite } X \wedge$

$Q \subseteq \text{carrier-vec } n \wedge \text{finite } Q \wedge$

$P = \text{convex-hull } Q + \text{cone } X \wedge$

$(A \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd \longrightarrow b \in \mathbb{Z}_v \cap \text{Bounded-vec } Bnd \longrightarrow$

$X \subseteq \mathbb{Z}_v \cap \text{Bounded-vec } (\text{det-bound } n \ (\text{max } 1 \ Bnd))$

$\wedge Q \subseteq \text{Bounded-vec } (\text{det-bound } n \ (\text{max } 1 \ Bnd)))$

<proof>

lemma *decomposition-theorem-polyhedra-2:*

assumes $Q: Q \subseteq \text{carrier-vec } n$ **and** $\text{fin-}Q: \text{finite } Q$

and $X: X \subseteq \text{carrier-vec } n$ **and** $\text{fin-}X: \text{finite } X$

and $P: P = \text{convex-hull } Q + \text{cone } X$

shows $\exists A \ b \ nr. A \in \text{carrier-mat } nr \ n \wedge b \in \text{carrier-vec } nr \wedge P = \text{polyhedron } A \ b$

<proof>

lemma *decomposition-theorem-polyhedra:*

$(\exists A \ b \ nr. A \in \text{carrier-mat } nr \ n \wedge b \in \text{carrier-vec } nr \wedge P = \text{polyhedron } A \ b)$

\longleftrightarrow

$(\exists Q \ X. Q \cup X \subseteq \text{carrier-vec } n \wedge \text{finite } (Q \cup X) \wedge P = \text{convex-hull } Q + \text{cone } X)$ **(is ?l = ?r)**

<proof>

lemma *polytope-equiv-bounded-polyhedron:*

polytope $P \longleftrightarrow$

$(\exists A \ b \ nr \ bnd. A \in \text{carrier-mat } nr \ n \wedge b \in \text{carrier-vec } nr \wedge P = \text{polyhedron } A \ b \wedge P \subseteq \text{Bounded-vec } bnd)$

<proof>

end

end

15 Mixed Integer Solutions

We prove that if an integral system of linear inequalities $Ax \leq b \wedge A'x < b'$ has a (mixed)integer solution, there there is also a small (mixed)integer solution, where the numbers are bounded by $(n + 1)! \cdot m^n$ where n is the number of variables and m is a bound on the absolute values of numbers occurring in A, A', b, b' .

theory *Mixed-Integer-Solutions*
imports *Decomposition-Theorem*
begin

definition *less-vec* :: 'a vec \Rightarrow ('a :: ord) vec \Rightarrow bool (**infix** $<_v$ 50) **where**
 $v <_v w = (\dim\text{-vec } v = \dim\text{-vec } w \wedge (\forall i < \dim\text{-vec } w. v \$ i < w \$ i))$

lemma *less-vecD*: **assumes** $v <_v w$ **and** $w \in \text{carrier-vec } n$
shows $i < n \Longrightarrow v \$ i < w \$ i$
<proof>

lemma *less-vecI*: **assumes** $v \in \text{carrier-vec } n$ $w \in \text{carrier-vec } n$
 $\bigwedge i. i < n \Longrightarrow v \$ i < w \$ i$
shows $v <_v w$
<proof>

lemma *less-vec-lesseq-vec*: $v <_v (w :: 'a :: \text{preorder vec}) \Longrightarrow v \leq w$
<proof>

lemma *floor-less*: $x \notin \mathbf{Z} \Longrightarrow \text{of-int } \lfloor x \rfloor < x$
<proof>

lemma *floor-of-int-eq[simp]*: $x \in \mathbf{Z} \Longrightarrow \text{of-int } \lfloor x \rfloor = x$
<proof>

locale *gram-schmidt-floor* = *gram-schmidt* n *f-ty*
for $n :: \text{nat}$ **and** *f-ty* :: 'a :: {*floor-ceiling*,
trivial-conjugatable-linordered-field} *itself*
begin

lemma *small-mixed-integer-solution-main*: **fixes** $A_1 :: 'a \text{ mat}$
assumes $A1: A_1 \in \text{carrier-mat } nr_1 \ n$
and $A2: A_2 \in \text{carrier-mat } nr_2 \ n$
and $b1: b_1 \in \text{carrier-vec } nr_1$
and $b2: b_2 \in \text{carrier-vec } nr_2$

and $A1Bnd$: $A_1 \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd$
and $b1Bnd$: $b_1 \in \mathbb{Z}_v \cap \text{Bounded-vec } Bnd$
and $A2Bnd$: $A_2 \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd$
and $b2Bnd$: $b_2 \in \mathbb{Z}_v \cap \text{Bounded-vec } Bnd$
and x : $x \in \text{carrier-vec } n$
and xI : $x \in \text{indexed-Ints-vec } I$
and sol-nonstrict : $A_1 *_v x \leq b_1$
and sol-strict : $A_2 *_v x <_v b_2$
shows $\exists x$.
 $x \in \text{carrier-vec } n \wedge$
 $x \in \text{indexed-Ints-vec } I \wedge$
 $A_1 *_v x \leq b_1 \wedge$
 $A_2 *_v x <_v b_2 \wedge$
 $x \in \text{Bounded-vec } (\text{fact } (n+1) * (\text{max } 1 \text{ Bnd}) \hat{\ } n)$
 (proof)

We get rid of the max-1 operation, by showing that a smaller value of Bnd can only occur in very special cases where the theorem is trivially satisfied.

lemma *small-mixed-integer-solution*: **fixes** $A_1 :: 'a \text{ mat}$
assumes $A1$: $A_1 \in \text{carrier-mat } nr_1 \ n$
and $A2$: $A_2 \in \text{carrier-mat } nr_2 \ n$
and $b1$: $b_1 \in \text{carrier-vec } nr_1$
and $b2$: $b_2 \in \text{carrier-vec } nr_2$
and $A1Bnd$: $A_1 \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd$
and $b1Bnd$: $b_1 \in \mathbb{Z}_v \cap \text{Bounded-vec } Bnd$
and $A2Bnd$: $A_2 \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd$
and $b2Bnd$: $b_2 \in \mathbb{Z}_v \cap \text{Bounded-vec } Bnd$
and x : $x \in \text{carrier-vec } n$
and xI : $x \in \text{indexed-Ints-vec } I$
and sol-nonstrict : $A_1 *_v x \leq b_1$
and sol-strict : $A_2 *_v x <_v b_2$
and non-degenerate : $nr_1 \neq 0 \vee nr_2 \neq 0 \vee Bnd \geq 0$
shows $\exists x$.
 $x \in \text{carrier-vec } n \wedge$
 $x \in \text{indexed-Ints-vec } I \wedge$
 $A_1 *_v x \leq b_1 \wedge$
 $A_2 *_v x <_v b_2 \wedge$
 $x \in \text{Bounded-vec } (\text{fact } (n+1) * Bnd \hat{\ } n)$
 (proof)

corollary *small-integer-solution-nonstrict*: **fixes** $A :: 'a \text{ mat}$
assumes A : $A \in \text{carrier-mat } nr \ n$
and b : $b \in \text{carrier-vec } nr$
and $ABnd$: $A \in \mathbb{Z}_m \cap \text{Bounded-mat } Bnd$
and $bBnd$: $b \in \mathbb{Z}_v \cap \text{Bounded-vec } Bnd$
and x : $x \in \text{carrier-vec } n$
and xI : $x \in \mathbb{Z}_v$
and sol : $A *_v x \leq b$

```

    and non-degenerate: nr ≠ 0 ∨ Bnd ≥ 0
  shows ∃ y.
    y ∈ carrier-vec n ∧
    y ∈ ℤv ∧
    A *v y ≤ b ∧
    y ∈ Bounded-vec (fact (n+1) * Bnd ^ n)
⟨proof⟩

end
end

```

16 Integer Hull

We define the integer hull of a polyhedron, i.e., the convex hull of all integer solutions. Moreover, we prove the result of Meyer that the integer hull of a polyhedron defined by an integer matrix is again a polyhedron, and give bounds for a corresponding decomposition theorem.

theory *Integer-Hull*

imports

Decomposition-Theorem

Mixed-Integer-Solutions

begin

context *gram-schmidt*

begin

definition *integer-hull* $P = \text{convex-hull } (P \cap \mathbb{Z}_v)$

lemma *integer-hull-mono*: $P \subseteq Q \implies \text{integer-hull } P \subseteq \text{integer-hull } Q$

⟨proof⟩

end

lemma *abs-neg-floor*: $|of-int\ b| \leq Bnd \implies -(\text{floor } Bnd) \leq b$

⟨proof⟩

lemma *abs-pos-floor*: $|of-int\ b| \leq Bnd \implies b \leq \text{floor } Bnd$

⟨proof⟩

context *gram-schmidt-floor*

begin

lemma *integer-hull-integer-cone*: **assumes** $C: C \subseteq \text{carrier-vec } n$

and $CI: C \subseteq \mathbb{Z}_v$

shows $\text{integer-hull } (\text{cone } C) = \text{cone } C$

⟨proof⟩

theorem *decomposition-theorem-integer-hull-of-polyhedron*: **assumes** $A: A \in \text{carrier-mat } nr\ n$

and b : $b \in \text{carrier-vec } nr$
and AI : $A \in \mathbb{Z}_m$
and bI : $b \in \mathbb{Z}_v$
and P : $P = \text{polyhedron } A \ b$
and Bnd : $Bnd = \text{Max } (\text{abs } ' (\text{elements-mat } A \cup \text{vec-set } b))$
shows $\exists H \ C. H \cup C \subseteq \text{carrier-vec } n \cap \mathbb{Z}_v$
 $\wedge H \subseteq \text{Bounded-vec } (\text{fact } (n+1) * (\text{max } 1 \ Bnd) \hat{\ }n)$
 $\wedge C \subseteq \text{Bounded-vec } (\text{fact } n * (\text{max } 1 \ Bnd) \hat{\ }n)$
 $\wedge \text{finite } (H \cup C)$
 $\wedge \text{integer-hull } P = \text{convex-hull } H + \text{cone } C$
 $\langle \text{proof} \rangle$

corollary integer-hull-of-polyhedron: assumes A : $A \in \text{carrier-mat } nr \ n$
and b : $b \in \text{carrier-vec } nr$
and AI : $A \in \mathbb{Z}_m$
and bI : $b \in \mathbb{Z}_v$
and P : $P = \text{polyhedron } A \ b$
shows $\exists A' \ b' \ nr'. A' \in \text{carrier-mat } nr' \ n \wedge b' \in \text{carrier-vec } nr' \wedge \text{integer-hull } P$
 $= \text{polyhedron } A' \ b'$
 $\langle \text{proof} \rangle$

corollary small-integer-solution-nonstrict-via-decomp: fixes $A :: 'a \ \text{mat}$
assumes A : $A \in \text{carrier-mat } nr \ n$
and b : $b \in \text{carrier-vec } nr$
and AI : $A \in \mathbb{Z}_m$
and bI : $b \in \mathbb{Z}_v$
and Bnd : $Bnd = \text{Max } (\text{abs } ' (\text{elements-mat } A \cup \text{vec-set } b))$
and x : $x \in \text{carrier-vec } n$
and xI : $x \in \mathbb{Z}_v$
and sol : $A * _v x \leq b$
shows $\exists y.$
 $y \in \text{carrier-vec } n \wedge$
 $y \in \mathbb{Z}_v \wedge$
 $A * _v y \leq b \wedge$
 $y \in \text{Bounded-vec } (\text{fact } (n+1) * (\text{max } 1 \ Bnd) \hat{\ }n)$
 $\langle \text{proof} \rangle$

end
end

References

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