

Lie Groups and Algebras

Richard Schmoetten and Jacques D. Fleuriot

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Abstract

Lie Groups are formalised as locales, building on the theory of Smooth Manifolds [1]. We formalise the diffeomorphism group of a manifold, and the action of a Lie group on a manifold. The general linear group is shown to be a Lie group by proving properties of the determinant, and matrix inverses. We also develop a theory of smooth vector fields on a C^∞ manifold M , defined as smooth maps from the manifold to its tangent bundle TM . We employ a shortcut that avoids difficulties in defining the tangent bundle as a manifold, but which still leads to vector fields with the properties one would expect. Notably, they are derivations $C^\infty(M) \rightarrow C^\infty(M)$. We construct *the* Lie algebra of a Lie group as an algebra of left-invariant smooth vector fields. Our main reference for the mathematics of smooth manifolds is Lee's textbook [2], which also contains material on Lie groups and algebras.

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theory *Algebra-On*

```

imports
  HOL-Types-To-Sets.Linear-Algebra-On
  Jacobson-Basic-Algebra.Ring-Theory
begin

```

1 Abstract algebra locales over a *field*

... with carrier set and some implicit operations (only algebraic multiplication, scaling, and derived constants are not implicit).

For full generality, one could define an algebra as a ring that is also a module (rather than a vector space, i.e. have a (non/commutative) base ring instead of a base field).

1.1 Bilinearity, Jacobi identity

```

lemma (in module-hom-on) mem-hom:
  assumes  $x \in S1$ 
  shows  $f x \in S2$ 
  using scale[OF assms, of 1] m2.mem-scale[of f x 1] m2.scale-one-on[of f x] oops

```

```

locale bilinear-on =
  vector-space-pair-on  $V W$  scaleV scaleW +
  vector-space-on  $X$  scaleX
  for  $V::'b::ab-group-add$  set and  $W::'c::ab-group-add$  set and  $X::'d::ab-group-add$ 
  set
  and  $scaleV::'a::field \Rightarrow 'b \Rightarrow 'b$  (infixr  $\bullet_V$  75)
  and  $scaleW::'a \Rightarrow 'c \Rightarrow 'c$  (infixr  $\bullet_W$  75)
  and  $scaleX::'a \Rightarrow 'd \Rightarrow 'd$  (infixr  $\bullet_X$  75) +
  fixes  $f::'b \Rightarrow 'c \Rightarrow 'd$ 
  assumes linearL:  $w \in W \Rightarrow linear-on\ V\ X\ scaleV\ scaleX\ (\lambda v. f\ v\ w)$ 
  and linearR:  $v \in V \Rightarrow linear-on\ W\ X\ scaleW\ scaleX\ (\lambda w. f\ v\ w)$ 
begin

```

```

lemma linearL':  $\llbracket v \in V; w \in W \rrbracket \Rightarrow f\ (a \bullet_V\ v)\ w = a \bullet_X\ (f\ v\ w)$ 
   $\llbracket v \in V; v' \in V; w \in W \rrbracket \Rightarrow f\ (v + v')\ w = (f\ v\ w) + (f\ v'\ w)$ 
  using linearL unfolding linear-on-def module-hom-on-def module-hom-on-axioms-def
  by simp+

```

```

lemma linearR':  $\llbracket v \in V; w \in W \rrbracket \Rightarrow f\ v\ (a \bullet_W\ w) = a \bullet_X\ (f\ v\ w)$ 
   $\llbracket v \in V; w \in W; w' \in W \rrbracket \Rightarrow f\ v\ (w + w') = (f\ v\ w) + (f\ v\ w')$ 
  using linearR unfolding linear-on-def module-hom-on-def module-hom-on-axioms-def
  by simp+

```

```

lemma bilinear-zero [simp]:
  shows  $w \in W \Rightarrow f\ 0\ w = 0$   $v \in V \Rightarrow f\ v\ 0 = 0$ 
  using linearL'(2) m1.mem-zero linearR'(2) m2.mem-zero by fastforce+

```

```

lemma bilinear-uminus [simp]:
  assumes  $v: v \in V$  and  $w: w \in W$ 
  shows  $f (-v) w = - (f v w)$   $f v (-w) = - (f v w)$ 
  using  $v w$  linearL'(2) m1.mem-uminus bilinear-zero(1) ab-left-minus add-right-imp-eq
apply metis
  using  $v w$  linearR'(2) m2.mem-uminus bilinear-zero(2) add-left-cancel add.right-inverse
by metis

```

end

For bilinear maps, "alternating" means the same as "skew-symmetric", which is the same as "anti-symmetric".

```

locale alternating-bilinear-on = bilinear-on  $S$   $S$   $S$  scale scale scale  $f$  for  $S$  scale  $f$ 
+
  assumes alternating:  $x \in S \implies f x x = 0$ 
begin

```

```

lemma antisym:
  assumes  $x \in S$   $y \in S$ 
  shows  $(f x y) + (f y x) = 0$ 
proof -
  have  $f (x+y) (x+y) = (f x x) + (f x y) + (f y x) + (f y y)$ 
    using linearL'(2) linearR'(2) by (simp add: assms m1.mem-add)
  thus ?thesis
    using alternating by (simp add: assms m1.mem-add)
qed

```

```

lemma antisym':
  assumes  $x \in S$   $y \in S$ 
  shows  $(f x y) = - (f y x)$ 
  using antisym[OF assms] by (simp add: eq-neg-iff-add-eq-0)

```

```

lemma antisym-uminus:
  assumes  $x \in S$   $y \in S$ 
  shows  $f (-x) y = f y x$   $f x (-y) = f y x$ 
  using bilinear-uminus by (metis antisym' assms)+

```

end

```

abbreviation (input) jacobi-identity-with:: $'a \Rightarrow ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow 'a \Rightarrow 'a \Rightarrow 'a \Rightarrow bool$ 
  where jacobi-identity-with zero-add f-add f-mult  $x y z \equiv$ 
     $zero-add = f-add (f-add (f-mult x (f-mult y z)) (f-mult y (f-mult z x))) (f-mult z (f-mult x y))$ 

```

```

abbreviation (input) jacobi-identity:: $('a::\{monoid-add\}) \Rightarrow 'a \Rightarrow 'a \Rightarrow 'a$ 

```

\Rightarrow *bool*
where *jacobi-identity f-mult x y z* \equiv *jacobi-identity-with 0 (+) f-mult x y z*

lemma (*in module-hom-on*) *mapsto-zero*: $f\ 0 = 0$
using *add m1.mem-zero* **by** *fastforce*

lemma (*in module-hom-on*) *mapsto-uminus*: $a \in S1 \implies f\ (-a) = -\ f\ a$
by (*metis add m1.mem-uminus neg-eq-iff-add-eq-0 mapsto-zero*)

lemma (*in module-hom-on*) *mapsto-closed*: $a \in S1 \implies f\ a \in S2$
using *mapsto-zero add mapsto-uminus*
oops

1.2 Unital and associative algebras

locale *algebra-on = bilinear-on S S S scale scale scale amult*
for *S*
and *scale* $:: 'a::field \Rightarrow 'b::ab-group-add \Rightarrow 'b$ (**infixr** $\langle *_S \rangle$ 75)
and *amult* (**infixr** $\langle \bullet \rangle$ 74) $+$
assumes *amult-closed* [*simp*]: $a \in S \implies b \in S \implies amult\ a\ b \in S$
begin

lemma
shows *distR*: $\llbracket x \in S; y \in S; z \in S \rrbracket \implies (x+y) \bullet z = x \bullet z + y \bullet z$
and *distL*: $\llbracket x \in S; y \in S; z \in S \rrbracket \implies z \bullet (x+y) = z \bullet x + z \bullet y$
and *scalar-compat* : $\llbracket x \in S; y \in S \rrbracket \implies (a *_S x) \bullet (b *_S y) = (a*b) *_S (x \bullet y)$
using *algebra-on-axioms unfolding algebra-on-def bilinear-on-def bilinear-on-axioms-def*
linear-on-def module-hom-on-def module-hom-on-axioms-def
by (*blast, blast, metis m1.mem-scale m1.scale-scale-on*)

lemma *scalar-compat'* [*simp*]:
shows $\llbracket x \in S; y \in S \rrbracket \implies (a *_S x) \bullet y = a *_S (x \bullet y)$
and $\llbracket x \in S; y \in S \rrbracket \implies x \bullet (a *_S y) = a *_S (x \bullet y)$
by (*simp-all add: linearL' linearR'*)

end

Sometimes an associative algebra is defined as a ring that is also a module (over a comm. ring), with the module and scalar multiplication being compatible, and the ring and module addition being the same. That definition implies an associative algebra is also unital, i.e. there is a multiplicative identity; in contrast, our definition doesn't. This is in agreement with how a *'a* needs no identity, and an additional type class *typ>'a::ring-1* is provided (instead of the terminology of *rng* vs. *ring*).

locale *assoc-algebra-on = algebra-on +*
assumes *amult-assoc*: $\llbracket x \in S; y \in S; z \in S \rrbracket \implies (x \bullet y) \bullet z = x \bullet (y \bullet z)$

locale *unital-algebra-on = algebra-on +*

```

fixes a-id
assumes amult-id [simp]:  $a-id \in S \implies a \in S \implies a \bullet a-id = a \implies a-id \bullet a = a$ 
begin

lemma id-neq-0-iff:  $\exists a \in S. \exists b \in S. a \neq b \iff 0 \neq a-id$ 
using amult-id(1) m1.mem-zero by blast

lemma id-neq-0-if:
shows  $a \in S \implies b \in S \implies a \neq b \implies 0 \neq a-id$ 
and  $card\ S \geq 2 \implies 0 \neq a-id$ 
and  $infinite\ S \implies 0 \neq a-id$ 
proof -

have ex-card:  $\exists S \subseteq A. card\ S = n$ 
if  $n \leq card\ A$ 
for  $n$  and  $A::'a\ set$ 
proof (cases finite A)
case True
from ex-bij-betw-nat-finite[OF this] obtain  $f$  where  $f: bij\ betw\ f\ \{0..<card\ A\}\ A\ ..$ 
moreover from  $f\ \langle n \leq card\ A \rangle$  have  $\{..<n\} \subseteq \{..<card\ A\}$  inj-on  $f\ \{..<n\}$ 
by (auto simp: bij-betw-def intro: subset-inj-on)
ultimately have  $f\ \{..<n\} \subseteq A$   $card\ (f\ \{..<n\}) = n$ 
by (auto simp: bij-betw-def card-image)
then show ?thesis by blast
next
case False
with  $\langle n \leq card\ A \rangle$  show ?thesis by force
qed

show  $a \in S \implies b \in S \implies a \neq b \implies 0 \neq a-id$ 
using amult-id(2) linearR' m1.mem-zero m1.scale-zero-left by metis
thus  $card\ S \geq 2 \implies 0 \neq a-id$ 
by (metis amult-id(2) card-2-iff' ex-card m1.mem-zero m1.scale-zero-left scalar-compat'(2) subset-iff)
thus  $infinite\ S \implies 0 \neq a-id$ 
using infinite-arbitrarily-large
by (metis amult-id(2) card-2-iff' linearR'(1) m1.mem-zero m1.scale-zero-left subset-iff)
qed

lemma id-neq-0-implies-elements :  $\exists a \in S. \exists b \in S. a \neq b$  if  $0 \neq a-id$ 
using amult-id(1) m1.mem-zero that by blast

lemma id-neq-0-implies-card:
assumes  $0 \neq a-id$ 
obtains  $card\ S \geq 2 \mid infinite\ S$ 

```

using *id-neq-0-implies-elements*[*OF assms*] **unfolding** *numeral-2-eq-2*
using *card-le-Suc0-iff-eq not-less-eq-eq* **by** *blast*

lemma *id-unique* [*simp*]:
fixes *other-id*
assumes *other-id* ∈ *S* ∧ *a*. *a* ∈ *S* ⇒ *a* • *other-id* = *a* ∧ *other-id* • *a* = *a*
shows *other-id* = *a-id*
using *assms amult-id* **by** *fastforce*

end

locale *assoc-algebra-1-on* = *assoc-algebra-on* + *unital-algebra-on* +
assumes *id-neq-0* [*simp*]: *a-id* ≠ 0 — this is as in the class *ring-1*, and merely
assures *S* has at least two elements
begin

lemma *is-ring-1-axioms*:
shows ∧ *a b c*. *a* ∈ *S* ⇒ *b* ∈ *S* ⇒ *c* ∈ *S* ⇒ *a* • *b* • *c* = *a* • (*b* • *c*)
and ∧ *a*. *a* ∈ *S* ⇒ *a-id* • *a* = *a*
and ∧ *a*. *a* ∈ *S* ⇒ *a* • *a-id* = *a*
and ∧ *a b c*. *a* ∈ *S* ⇒ *b* ∈ *S* ⇒ *c* ∈ *S* ⇒ (*a* + *b*) • *c* = *a* • *c* + *b* • *c*
and ∧ *a b c*. *a* ∈ *S* ⇒ *b* ∈ *S* ⇒ *c* ∈ *S* ⇒ *a* • (*b* + *c*) = *a* • *b* + *a* • *c*
by (*simp-all add: distR distL algebra-simps*)

lemma *inverse-unique* [*simp*]:
assumes *a*: *a* ∈ *S* *a* ≠ 0
and *x*: *x* ∈ *S* *a* • *x* = *a-id* ∧ *x* • *a* = *a-id*
and *y*: *y* ∈ *S* *a* • *y* = *a-id* ∧ *y* • *a* = *a-id*
shows *x* = *y*
using *amult-assoc*[*of x a x*] *amult-assoc*[*of x a y*]
by (*simp add: assms*)

lemma *inverse-unique'*:
assumes *a*: *a* ∈ *S* *a* ≠ 0
and *inv-ex*: ∃ *x* ∈ *S*. *a* • *x* = *a-id* ∧ *x* • *a* = *a-id*
shows ∃! *x* ∈ *S*. *a* • *x* = *a-id* ∧ *x* • *a* = *a-id*
using *a inv-ex inverse-unique* **by** (*metis (no-types, lifting)*)

end

lemma *algebra-onI* [*intro*]:
fixes *scale* :: '*a*::*field* ⇒ '*b*::*ab-group-add* ⇒ '*b* (**infixr** *_{*S*} 75)
and *amult* (**infixr** • 74)
assumes *vector-space-on S scale*
and *distR*: ∧ *x y z*. [[*x* ∈ *S*; *y* ∈ *S*; *z* ∈ *S*]] ⇒ (*x* + *y*) • *z* = *x* • *z* + *y* • *z*
and *distL*: ∧ *x y z*. [[*x* ∈ *S*; *y* ∈ *S*; *z* ∈ *S*]] ⇒ *z* • (*x* + *y*) = *z* • *x* + *z* • *y*
and *scalar-compat*: ∧ *a x y*. [[*x* ∈ *S*; *y* ∈ *S*]] ⇒ (*a* *_{*S*} *x*) • *y* = *a* *_{*S*} (*x* • *y*) ∧ *x* •
(*a* *_{*S*} *y*) = *a* *_{*S*} (*x* • *y*)

and closure: $\bigwedge x y. \llbracket x \in S; y \in S \rrbracket \implies x \bullet y \in S$
shows algebra-on S scale amult
unfolding algebra-on-def bilinear-on-def vector-space-pair-on-def bilinear-on-axioms-def
apply (*intro conjI algebra-on-axioms.intro, simp-all add: assms(1)*)
unfolding linear-on-def module-hom-on-def module-hom-on-axioms-def
by (*auto simp: assms vector-space-on.axioms*)

lemma (*in vector-space-on*) *scalar-compat-iff:*

fixes *scale-notation* (**infixr** $*_S$ 75)
and *amult* (**infixr** \bullet 74)
defines *scale-notation* \equiv *scale*
assumes *distR:* $\bigwedge x y z. \llbracket x \in S; y \in S; z \in S \rrbracket \implies (x+y) \bullet z = x \bullet z + y \bullet z$
and *distL:* $\bigwedge x y z. \llbracket x \in S; y \in S; z \in S \rrbracket \implies z \bullet (x+y) = z \bullet x + z \bullet y$
shows $(\forall a. \forall x \in S. \forall y \in S. (a *_S x) \bullet y = a *_S (x \bullet y) \wedge x \bullet (a *_S y) = a *_S (x \bullet y)) \iff$
 $(\forall a b. \forall x \in S. \forall y \in S. (a *_S x) \bullet (b *_S y) = (a *_S b) *_S (x \bullet y))$
proof (*intro iffI*)
{ **assume** *asm:* $\bigwedge a b x y. x \in S \implies y \in S \implies a *_S x \bullet b *_S y = (a *_S b) *_S (x \bullet y)$
}
{ **fix** *a x y*
assume *S:* $x \in S y \in S$
have $a *_S x \bullet y = a *_S (x \bullet y) \wedge x \bullet a *_S y = a *_S (x \bullet y)$
using *asm[of x y a 1] S* **apply** (*simp add: scale-notation-def*)
using *asm[of x y 1 a] S* **by** (*simp add: scale-notation-def*) **}}**
thus $\forall a b. \forall x \in S. \forall y \in S. a *_S x \bullet b *_S y = (a *_S b) *_S (x \bullet y) \implies$
 $\forall a. \forall x \in S. \forall y \in S. a *_S x \bullet y = a *_S (x \bullet y) \wedge x \bullet a *_S y = a *_S (x \bullet y)$
by *blast*
qed (*metis mem-scale scale-notation-def scale-scale-on*)

lemma (*in vector-space-on*) *algebra-onI:*

fixes *scale-notation* (**infixr** $*_S$ 75)
and *amult* (**infixr** \bullet 74)
defines *scale-notation* \equiv *scale*
assumes *distR:* $\bigwedge x y z. \llbracket x \in S; y \in S; z \in S \rrbracket \implies (x+y) \bullet z = x \bullet z + y \bullet z$
and *distL:* $\bigwedge x y z. \llbracket x \in S; y \in S; z \in S \rrbracket \implies z \bullet (x+y) = z \bullet x + z \bullet y$
and *scalar-compat:* $\bigwedge a x y. \llbracket x \in S; y \in S \rrbracket \implies (a *_S x) \bullet y = a *_S (x \bullet y) \wedge x \bullet (a *_S y) = a *_S (x \bullet y)$
and *closure:* $\bigwedge x y. \llbracket x \in S; y \in S \rrbracket \implies x \bullet y \in S$
shows algebra-on S scale amult
using *algebra-onI[of S scale amult] assms scale-notation-def vector-space-on-axioms*
by *simp*

1.3 Lie algebra (locale)

List syntax interferes with the standard notation for the Lie bracket, so it can be disabled it here. Instead, we add a delimiter to the notation for Lie brackets, which also helps with unambiguous parsing.

locale *lie-algebra* = *algebra-on* \mathfrak{g} *scale* *lie-bracket* + *alternating-bilinear-on* \mathfrak{g} *scale* *lie-bracket*

for \mathfrak{g}
and *scale* :: 'a::field \Rightarrow 'b::ab-group-add \Rightarrow 'b (**infixr** $\langle *_S \rangle$ 75)
and *lie-bracket* :: 'b \Rightarrow 'b \Rightarrow 'b ($\langle [-;-] \rangle$ 74) +
assumes *jacobi*: $\llbracket x \in \mathfrak{g}; y \in \mathfrak{g}; z \in \mathfrak{g} \rrbracket \Longrightarrow 0 = [x; [y; z]] + [y; [z; x]] + [z; [x; y]]$

lemma (**in** *algebra-on*) *lie-algebraI*:

assumes *alternating*: $\forall x \in S. \text{amult } x \ x = 0$
and *jacobi*: $\forall x \in S. \forall y \in S. \forall z \in S. \text{jacobi-identity } \text{amult } x \ y \ z$
shows *lie-algebra* S *scale* *amult*
apply *unfold-locales* **using** *assms* **by** *auto*

lemma (**in** *vector-space-on*) *lie-algebraI*:

fixes *lie-bracket* :: 'b \Rightarrow 'b \Rightarrow 'b ($\langle [-;-] \rangle$ 74)
and *scale-notation* (**infixr** $*_S$ 75)
defines *scale-notation* \equiv *scale*
assumes *distributivity*:
 $\bigwedge x \ y \ z. \llbracket x \in S; y \in S; z \in S \rrbracket \Longrightarrow [(x+y); z] = [x; z] + [y; z] \wedge [z; (x+y)] = [z; x] + [z; y]$
and *scalar-compatibility*:
 $\bigwedge a \ x \ y. \llbracket x \in S; y \in S \rrbracket \Longrightarrow [(a *_S x); y] = a *_S ([x; y]) \wedge [x; (a *_S y)] = a *_S ([x; y])$
and *closure*: $\bigwedge x \ y. \llbracket x \in S; y \in S \rrbracket \Longrightarrow [x; y] \in S$
and *alternating*: $\forall x \in S. \text{lie-bracket } x \ x = 0$
and *jacobi*: $\forall x \in S. \forall y \in S. \forall z \in S. \text{jacobi-identity } \text{lie-bracket } x \ y \ z$
shows *lie-algebra* S *scale* *lie-bracket*
using *assms*(1,3,6) **by** (*auto simp: assms*(2,4,5) *intro!: algebra-on.lie-algebraI algebra-onI*)

context *lie-algebra* **begin**

lemma *jacobi-alt*:

assumes $x: x \in \mathfrak{g}$ **and** $y: y \in \mathfrak{g}$ **and** $z: z \in \mathfrak{g}$
shows $[x; [y; z]] = [[x; y]; z] + [y; [x; z]]$
proof –
have $[x; [y; z]] = - ([y; [z; x]]) + (- ([z; [x; y]]))$
using *jacobi*[*OF assms*] *add-implies-diff*[*of* $[x; [y; z]]$ $[y; [z; x]] + [z; [x; y]]$ 0]
by (*simp add: add commute add.left-commute*)
moreover **have** $[[x; y]; z] = - ([z; [x; y]]) - ([y; [z; x]]) = [y; [x; z]]$
using *antisym'*[*OF amult-closed*[*OF* $x \ y$] z] *antisym'*[*OF* $z \ x$] **by** (*simp-all add: assms*)
ultimately show $[x; [y; z]] = [[x; y]; z] + [y; [x; z]]$ **by** *simp*
qed

lemma *lie-subalgebra*:

assumes $h: \mathfrak{h} \subseteq \mathfrak{g}$ *m1.subspace* \mathfrak{h} **and** *closed*: $\bigwedge x \ y. x \in \mathfrak{h} \Longrightarrow y \in \mathfrak{h} \Longrightarrow \text{lie-bracket } x \ y \in \mathfrak{h}$

```

shows lie-algebra  $\mathfrak{h}$  scale lie-bracket
proof –
interpret  $\mathfrak{h}$ : vector-space-on  $\mathfrak{h}$  scale
  apply unfold-locales
  apply (meson h(1) m1.scale-right-distrib-on subset-iff)
  apply (meson h(1) in-mono m1.scale-left-distrib-on)
  using h(1) m1.scale-scale-on m1.scale-one-on apply auto[2]
  by (simp-all add: h m1.subspace-add m1.subspace-0 m1.subspace-scale)

show ?thesis
  apply (intro  $\mathfrak{h}$ .lie-algebraI)
  using alternating h(1) jacobi linearL' linearR' by (auto simp: closed subset-iff)
qed

end

```

1.4 Division algebras

abbreviation (in algebra-on) *is-left-divisor* $x a b \equiv x \in S \wedge a = \text{amult } x b$
abbreviation (in algebra-on) *is-right-divisor* $x a b \equiv x \in S \wedge a = \text{amult } b x$

```

locale div-algebra-on = algebra-on +
  fixes divL::'a $\Rightarrow$ 'a $\Rightarrow$ 'a
  and divR::'a $\Rightarrow$ 'a $\Rightarrow$ 'a
  assumes divL:  $\llbracket a \in S; b \in S; b \neq 0 \rrbracket \implies \text{is-left-divisor } (\text{divL } a b) a b$ 
   $\llbracket a \in S; b \in S; b \neq 0 \rrbracket \implies \text{is-left-divisor } y a b \implies y = (\text{divL } a b)$ 
  and divR:  $\llbracket a \in S; b \in S; b \neq 0 \rrbracket \implies \text{is-right-divisor } (\text{divR } a b) a b$ 
   $\llbracket a \in S; b \in S; b \neq 0 \rrbracket \implies \text{is-right-divisor } y a b \implies y = (\text{divR } a b)$ 
begin

```

In terms of the vocabulary of division rings, the expression $a = \text{divL } a b$ $b \bullet b$ means that $\text{divL } a b$ is a left divisor of a , and conversely that a is a right multiple of $\text{divL } a b$.

For $b = (0::'c)$, the divisors still exist as members of the correct type (necessarily), but they have no properties. Similarly for correctly-typed input outside the algebra.

```

lemma [simp]:
  assumes  $a \in S \ b \in S \ b \neq 0$ 
  shows divL':  $\text{divL } a b \in S \ (\text{divL } a b) \bullet b = a \ \forall y \in S. a = y \bullet b \longrightarrow y = \text{divL } a b$ 
  and divR':  $\text{divR } a b \in S \ b \bullet (\text{divR } a b) = a \ \forall y \in S. a = b \bullet y \longrightarrow y = \text{divR } a b$ 
  using assms divL divR by simp-all
end

```

```

lemma (in algebra-on) div-algebra-onI:
  assumes  $\forall a \in S. \forall b \in S. b \neq 0 \longrightarrow (\exists! x \in S. a = b \bullet x) \wedge (\exists! y \in S. a = y \bullet b)$ 
  shows div-algebra-on S scale amult ( $\lambda a b. \text{THE } y. y \in S \wedge a = y \bullet b$ ) ( $\lambda a b. \text{THE } x. x \in S \wedge a = b \bullet x$ )

```

proof (*unfold div-algebra-on-def div-algebra-on-axioms-def, intro conjI allI impI*)
fix $a\ b\ x$
assume $a \in S\ b \in S\ b \neq 0$
have $exL: \exists!x \in S. a = x \bullet b$ **by** (*simp add: $\langle a \in S \rangle \langle b \in S \rangle \langle b \neq 0 \rangle$ assms*)
from $theI'[OF\ this]$
show $L: (THE\ y. y \in S \wedge a = y \bullet b) \in S\ a = (THE\ y. y \in S \wedge a = y \bullet b) \bullet b$
by *simp+*
have $exR: \exists!x \in S. a = b \bullet x$ **by** (*simp add: $\langle a \in S \rangle \langle b \in S \rangle \langle b \neq 0 \rangle$ assms*)
from $theI'[OF\ this]$
show $R: (THE\ x. x \in S \wedge a = b \bullet x) \in S\ a = b \bullet (THE\ x. x \in S \wedge a = b \bullet x)$
by *simp+*
{ **assume** $x \in S \wedge a = x \bullet b$
thus $x = (THE\ y. y \in S \wedge a = y \bullet b)$ **using** $L\ exL$ **by** *auto*
} { **assume** $x \in S \wedge a = b \bullet x$
thus $x = (THE\ x. x \in S \wedge a = b \bullet x)$ **using** $R\ exR$ **by** *auto*
}
qed (*simp add: algebra-on-axioms*)

lemma (**in** *assoc-algebra-1-on*) *div-algebra-onI'*:

fixes $ainv\ adivL\ adivR$
defines $ainv\ a \equiv (THE\ x. x \in S \wedge a-id = x \bullet a \wedge a-id = a \bullet x)$
and $adivL\ b\ a \equiv b \bullet (ainv\ a)$
and $adivR\ b\ a \equiv (ainv\ a) \bullet b$
assumes $\forall a \in S. a \neq 0 \longrightarrow (\exists x \in S. a-id = x \bullet a \wedge a-id = a \bullet x)$
shows *div-algebra-on S scale amult adivL adivR*

proof (*unfold-locales*)

fix $a\ b$
assume $asm: a \in S\ b \in S\ b \neq 0$
have $inv-ex: \exists!x \in S. a-id = x \bullet b \wedge a-id = b \bullet x$
using $assms(4)\ inverse-unique'\ asm(2,3)$ **by** *metis*
let $?a = THE\ x. x \in S \wedge a-id = x \bullet b \wedge a-id = b \bullet x$
from $theI'[OF\ inv-ex]$ **show** $1: adivR\ a\ b \in S \wedge a = b \bullet adivR\ a\ b$
unfolding *adivR-def ainv-def* **apply** (*intro conjI*)
using $asm(1)$ **apply** *simp*
using *amult-assoc amult-id(2) asm(1,2) is-ring-1-axioms(2)* **by** (*metis (no-types, lifting)*)
from $theI'[OF\ inv-ex]$ **show** $2: adivL\ a\ b \in S \wedge a = adivL\ a\ b \bullet b$
unfolding *adivL-def ainv-def* **apply** (*intro conjI*)
apply (*simp add: asm(1)*)
using *amult-assoc asm(1,2) is-ring-1-axioms(3)* **by** *presburger*
{ **fix** y **assume** $y \in S \wedge a = y \bullet b$
thus $y = adivL\ a\ b$
by (*metis inv-ex 2 amult-assoc asm(2) amult-id(2)*)
} { **fix** y **assume** $y \in S \wedge a = b \bullet y$
thus $y = adivR\ a\ b$
by (*metis 1 amult-assoc asm(2) inv-ex is-ring-1-axioms(2)*) }
qed

lemma (**in** *assoc-algebra-on*) *div-algebra-on-imp-inverse*:

assumes *div-algebra-on S scale amult divL divR card S ≥ 2 ∨ infinite S*
shows $\exists a-id \in S. (\forall a \in S. a \bullet a-id = a \wedge a-id \bullet a = a) \wedge (\forall a \in S. a \neq 0 \longrightarrow \text{divL } a-id \ a = \text{divR } a-id \ a)$
proof –
obtain *x where x ≠ 0 x ∈ S*
using *assms(2) unfolding numeral-2-eq-2*
by (*metis card-1-singleton-iff card-gt-0-iff card-le-Suc0-iff-eq insertI1 not-less-eq-eq rev-finite-subset subsetI zero-less-Suc*)
let *?id = divL x x*
show *?thesis*
proof (*intro bexI conjI ballI impI*)
show *1: ?id ∈ S*
using *assms unfolding div-algebra-on-def div-algebra-on-axioms-def*
using $\langle x \in S \rangle \langle x \neq 0 \rangle$ **by** *blast*
fix *a assume a ∈ S*
show *2: a • ?id = a*
by (*smt (verit) 1 ⟨a ∈ S⟩ ⟨x ∈ S⟩ ⟨x ≠ 0⟩ amult-assoc amult-closed assms(1) div-algebra-on.divL*)
show *3: ?id • a = a*
by (*smt (verit) ⟨a ∈ S⟩ ⟨x ∈ S⟩ ⟨x ≠ 0⟩ amult-assoc assms(1) div-algebra-on.divL(1) div-algebra-on.divR'*)
assume *a ≠ 0*
show *4: divL ?id a = divR ?id a*
by (*smt (verit) 1 3 ⟨a ∈ S⟩ ⟨a ≠ 0⟩ amult-assoc amult-closed assms(1) div-algebra-on.divL div-algebra-on.divR(2)*)
qed
qed

lemma (*in assoc-algebra-on*) *assoc-div-algebra-on-iff*:
assumes *card S ≥ 2 ∨ infinite S*
shows $(\exists \text{divL divR. div-algebra-on } S \text{ scale amult divL divR}) \longleftrightarrow$
 $(\exists id. \text{unital-algebra-on } S \text{ scale amult id} \wedge (\forall a \in S. a \neq 0 \longrightarrow (\exists x \in S. a \bullet x = id \wedge x \bullet a = id)))$
proof (*intro iffI*)
assume $\exists id. \text{unital-algebra-on } S (*_S) (\bullet) id \wedge (\forall a \in S. a \neq 0 \longrightarrow (\exists x \in S. a \bullet x = id \wedge x \bullet a = id))$
then obtain *id*
where *id: id ∈ S* $\forall a \in S. a \bullet id = a \wedge id \bullet a = a$ **and** *inv: $\forall a \in S. a \neq 0 \longrightarrow (\exists x \in S. a \bullet x = id \wedge x \bullet a = id)$*
using *unital-algebra-on.amult-id by blast*
then have *unital: unital-algebra-on S scale amult id*
by (*unfold-locales, simp-all*)
then have *assoc-alg: assoc-algebra-1-on S scale amult id*
unfolding *assoc-algebra-1-on-def assoc-algebra-1-on-axioms-def*
using *assms unital-algebra-on.id-neq-0-if(2,3) assoc-algebra-on-axioms*
by *blast*
show $\exists \text{divL divR. div-algebra-on } S (*_S) (\bullet) \text{divL divR}$
using *assoc-algebra-1-on.div-algebra-onI'[OF assoc-alg] inv by fastforce*
next

```

assume  $\exists \text{divL divR. div-algebra-on } S (*_S) (\bullet) \text{divL divR}$ 
then obtain  $\text{divL divR}$  where  $\text{div-alg: div-algebra-on } S (*_S) (\bullet) \text{divL divR}$  by
blast
show  $\exists \text{id. unital-algebra-on } S (*_S) (\bullet) \text{id} \wedge (\forall a \in S. a \neq 0 \longrightarrow (\exists x \in S. a \bullet x = \text{id} \wedge x \bullet a = \text{id}))$ 
using  $\text{div-algebra-on-imp-inverse[OF div-alg assms] unital-algebra-on-axioms.intro assoc-algebra-on-axioms}$ 
unfolding  $\text{unital-algebra-on-def unital-algebra-on-axioms-def assoc-algebra-on-def}$ 
by  $(\text{smt (verit) div-alg div-algebra-on.divL(1) div-algebra-on.divR(1)})$ 
qed

```

```

locale  $\text{assoc-div-algebra-on} =$ 
 $\text{assoc-algebra-1-on } S \text{ scale amult a-id} +$ 
 $\text{div-algebra-on } S \text{ scale amult } \lambda a b. \text{amult } a (a\text{-inv } b) \lambda a b. \text{amult } (a\text{-inv } b) a$ 
for  $S$ 
and  $\text{scale} :: 'a::\text{field} \Rightarrow 'b::\text{ab-group-add} \Rightarrow 'b$  (infixr  $\langle *_S \rangle$  75)
and  $\text{amult} :: 'b \Rightarrow 'b \Rightarrow 'b$  (infixr  $\langle \bullet \rangle$  74)
and  $\text{a-id} :: 'b \langle \mathbf{1} \rangle$ 
and  $\text{a-inv} :: 'b \Rightarrow 'b$ 
begin

```

The definition *assoc-div-algebra-on* is justified by $2 \leq \text{card } S \vee \text{infinite } S \implies (\exists \text{divL divR. div-algebra-on } S (*_S) (\bullet) \text{divL divR}) = (\exists \text{id. unital-algebra-on } S (*_S) (\bullet) \text{id} \wedge (\forall a \in S. a \neq (0::'b) \longrightarrow (\exists x \in S. a \bullet x = \text{id} \wedge x \bullet a = \text{id})))$ above: If we have an associative algebra already, the only way it can be a division algebra is to be unital as well. Since now left and right divisors can be defined through multiplicative inverses, we take only the inverse as a locale parameter, and construct the divisors. The only case we miss here (due to the requirement $\mathbf{1} \neq (0::'b)$) is the trivial algebra, which contains only the zero element (which acts as identity as well). This is for compatibility with the standard Isabelle/HOL type classes, which are subclasses of *zero-neq-one*.

```

abbreviation  $(\text{input}) \text{divL} :: 'b \Rightarrow 'b \Rightarrow 'b$ 
where  $\text{divL } a b \equiv \text{amult } a (a\text{-inv } b)$ 

```

```

abbreviation  $(\text{input}) \text{divR} :: 'b \Rightarrow 'b \Rightarrow 'b$ 
where  $\text{divR } a b \equiv \text{amult } (a\text{-inv } b) a$ 

```

```

lemma  $\text{div-self-eq-id}$ :
assumes  $a \in S \ a \neq 0$ 
shows  $\text{divL } a a = a\text{-id}$ 
and  $\text{divR } a a = a\text{-id}$ 
apply  $(\text{metis amult-id}(1,3) \text{assms divL}'(3))$ 
by  $(\text{metis amult-id}(1,2) \text{assms divR}'(3))$ 

```

```

end

```

```

locale finite-dimensional-assoc-div-algebra-on =
  assoc-div-algebra-on S scale amult a-id a-inv +
  finite-dimensional-vector-space-on S scale basis
for S :: ⟨'b::ab-group-add set⟩
  and scale :: ⟨'a::field ⇒ 'b ⇒ 'b⟩ (infixr ⟨*_S⟩ 75)
  and amult :: ⟨'b⇒'b⇒'b⟩ (infixr ⟨•⟩ 74)
  and a-id :: ⟨'b⟩ (⟨1⟩)
  and a-inv :: ⟨'b⇒'b⟩
  and basis :: ⟨'b set⟩

lemma (in assoc-div-algebra-on) finite-dimensional-assoc-div-algebra-onI [intro]:
  fixes basis :: 'b set
  assumes finite-Basis: finite basis
  and independent-Basis: ¬ m1.dependent basis
  and span-Basis: m1.span basis = S
  and basis-subset: basis ⊆ S
  shows finite-dimensional-assoc-div-algebra-on S scale amult a-id a-inv basis
  by (unfold-locales, simp-all add: assms)

end

```

```

theory Linear-Algebra-More
imports
  HOL-Analysis.Analysis
  Smooth-Manifolds.Smooth
  Transfer-Cayley-Hamilton
begin

```

2 Continuity of the determinant (and other maps)

```

lemma continuous-on-proj: continuous-on s fst continuous-on s snd
  apply (simp add: continuous-on-fst[OF continuous-on-id])
  by (simp add: continuous-on-snd[OF continuous-on-id])

lemma continuous-on-plus:
  fixes s::('a × 'a::topological-monoid-add) set
  shows continuous-on s (λ(x,y). x+y)
  by (simp add: continuous-on-add[OF continuous-on-proj] case-prod-beta')

lemma continuous-on-times:
  fixes s::('a × 'a::real-normed-algebra) set
  shows continuous-on s (λ(x,y). x*y)
  by (simp add: case-prod-beta' continuous-on-mult[OF continuous-on-proj])

lemma continuous-on-times':
  fixes s::('a × 'a::topological-monoid-mult) set

```

shows *continuous-on* s ($\lambda(x,y). x*y$)
by (*simp add: case-prod-beta' continuous-on-mult'[OF continuous-on-proj]*)

Only functions between *real-normed-vector* spaces can be *bounded-linear*...

lemma *continuous-on-nth-of-vec*:
fixes $s::('a::\text{real-normed-field}, 'n::\text{finite})\text{vec set}$
shows *continuous-on* s ($\lambda x. x \$ n$)
by (*simp add: bounded-linear-vec-nth linear-continuous-on*)

lemma *bounded-linear-mat-ijth[intro]*: *bounded-linear* ($\lambda x. x \$ i \$ j$)
apply (*standard; simp?*)
apply (*intro exI[of - 1]*)
apply (*simp add: norm-nth-le*)
by (*meson Finite-Cartesian-Product.norm-nth-le dual-order.trans*)

lemma *continuous-on-ijth-of-mat*:
fixes $s::('a::\text{real-normed-field}, 'n::\text{finite})\text{square-matrix set}$
shows *continuous-on* s ($\lambda x. x \$ i \$ j$)
by (*simp add: bounded-linear-mat-ijth linear-continuous-on*)

lemma *continuous-on-det*:
fixes $s::('a::\text{real-normed-field}, 'n::\text{finite})\text{square-matrix set}$
shows *continuous-on* s *det*
proof (*unfold det-def, intro continuous-on-sum*)
fix p
assume $p \in \{p. p \text{ permutes } (UNIV::'n \text{ set})\}$
show *continuous-on* s ($\lambda A. \text{of-int } (\text{sign } p) * (\prod_{i \in UNIV}. A \$ i \$ p \ i)$)
proof (*intro continuous-on-mult*)
show *continuous-on* s ($\lambda x. \text{of-int } (\text{sign } p)$)
by *simp*
show *continuous-on* s ($\lambda x. \prod_{i \in UNIV}. x \$ i \$ p \ i$)
apply (*intro continuous-on-prod*)
by (*simp add: continuous-on-ijth-of-mat*)
qed
qed

lemma *invertible-inv-ex*:
fixes $a::'a::\text{semiring-1}^{\wedge}n^{\wedge}n$
assumes *invertible* a
shows $(\text{matrix-inv } a)**a = \text{mat } 1 \ a**(\text{matrix-inv } a) = \text{mat } 1$
using *some-eq-ex assms invertible-def matrix-inv-def*
by (*smt (verit, ccfv-SIG)+*)

A similar result to the below already exists for fields, see e.g. *invertible-left-inverse*. This is more general, as it applies to any semiring (with 1).

lemma *invertible-matrix-inv*:
fixes $a::'a::\text{semiring-1}^{\wedge}n^{\wedge}n$
assumes *invertible* a

shows *invertible (matrix-inv a)*
 using *invertible-inv-ex assms invertible-def*
 by *auto*

3 Component expressions for inverse matrices over fields

lemma *inv-adj-det-field-component*:
 fixes $i\ j::'n::\text{finite}$ and $A\ A'::'a::\text{field}^n$
 defines $\text{inv}A: A' \equiv \text{map-matrix } (\lambda x. x / (\text{det } A)) (\text{adjugate } A)$
 assumes *invertible A*
 shows $(A ** A') \$i \$j = (\text{if } i=j \text{ then } 1 \text{ else } 0)$
proof –
 let $?D = \text{det } A$
 have *det-not-0: ?D ≠ 0*
 using *assms by (metis det-I det-mul invertible-inv-ex(2) mult-zero-left zero-neq-one)*
 have $(\sum_{k \in \text{UNIV}} A \$i \$k * (\text{adjugate } A) \$k \$j) = (\text{if } i=j \text{ then } ?D \text{ else } 0)$
 using *mult-adjugate-det-2 [of A] unfolding matrix-matrix-mult-def mat-def*
 by *(metis (mono-tags, lifting) iso-tuple-UNIV-I vec-lambda-inverse)*
 then have $(\text{if } i=j \text{ then } 1 \text{ else } 0) = (\sum_{k \in \text{UNIV}} A \$i \$k * (\text{adjugate } A) \$k \$j) / ?D$
 by *(simp add: det-not-0)*
 also have $\dots = (\sum_{k \in \text{UNIV}} A \$i \$k * A' \$k \$j)$
 using *sum-divide-distrib invA by force*
 finally show *?thesis*
 unfolding *matrix-matrix-mult-def by simp*
qed

lemma *inverse-adjugate-det-2*:
 fixes $A::'a::\text{field}^n$
 assumes *invertible A*
 shows $\text{matrix-inv } A = \text{map-matrix } (\lambda x. x / (\text{det } A)) (\text{adjugate } A)$
 (*is matrix-inv A = ?A'*)
proof –
 let $?D = \text{det } A$
 have *det-not-0: ?D ≠ 0*
 using *assms by (metis det-I det-mul invertible-inv-ex(2) mult-zero-left zero-neq-one)*
 have $AA': A ** ?A' = \text{mat } 1$
 unfolding *mat-def using inv-adj-det-field-component[OF assms] by (simp add: vec-eq-iff)*
 moreover have $?A' ** A = \text{mat } 1$
 using *AA' by (simp add: matrix-left-right-inverse)*
 ultimately show $\text{matrix-inv } A = ?A'$
 by *(metis (no-types) invertible-def invertible-inv-ex(2) matrix-mul-assoc matrix-mul-lid)*
qed

lemma *inverse-adjugate-det*:

fixes $A::'a::\text{field}^{\wedge n} \wedge n$
assumes *invertible* A
shows $\text{matrix-inv } A = (1 / (\det A)) *_s (\text{adjugate } A)$
using *inverse-adjugate-det-2* [*OF assms*] **unfolding** *map-matrix-def smult-mat-def*
by *auto*

lemma *transpose-component*: $(\text{transpose } A) \$i\$j = A\$j\i
unfolding *transpose-def* **by** *simp*

lemma *matrix-inverse-component*:
fixes $A::'a::\text{field}^{\wedge n} \wedge n$ **and** $i\ j::'n::\text{finite}$
assumes *invertible* A
shows $(\text{matrix-inv } A) \$i\$j = \det (\chi\ k\ l. \text{ if } k = j \wedge l = i \text{ then } 1 \text{ else if } k = j \vee l = i \text{ then } 0 \text{ else } A \$k\$l) / (\det A)$
using *inverse-adjugate-det-2* [*OF assms*]
by (*simp add: transpose-component adjugate-def cofac-def minor-mat-def*)

lemma *matrix-adjugate-component*:
fixes $A::'a::\text{field}^{\wedge n} \wedge n$ **and** $i\ j::'n::\text{finite}$
assumes *invertible* A
shows $(\text{adjugate } A) \$i\$j = \det (\chi\ k\ l. \text{ if } k = j \wedge l = i \text{ then } 1 \text{ else if } k = j \vee l = i \text{ then } 0 \text{ else } A \$k\$l)$
by (*simp add: transpose-component adjugate-def cofac-def minor-mat-def*)

4 Smoothness of real matrix operations and *det*

4.1 Smoothness of matrix multiplication

lemma *smooth-on-ijth-of-mat*:
fixes $s::('a::\text{real-normed-field}, 'n::\text{finite})\text{square-matrix set}$
shows *smooth-on* $s (\lambda x. x \$i\$j)$
by (*simp add: bounded-linear.smooth-on bounded-linear-mat-ijth*)

Notice the following result holds only for matrices over the real numbers. (Try removing the type annotations: Isabelle automatically casts to the indicated type anyway.) This is because only real inner product spaces are defined: thus whatever "base field" a matrix is defined over, is implicitly assumed to also be a real inner product space (as is possible, for example, for \mathbb{C} with the normal inner product of \mathbb{R}^2), and the inner product is built on top of the existing one to return a *real* result.

lemma *matrix-matrix-mul-component-real*:
fixes $A::\text{real}^{\wedge k} \wedge n$
and $B::\text{real}^{\wedge m} \wedge k$
shows $A**B = (\chi\ i\ j. \text{ inner } (\text{row } i\ A) (\text{column } j\ B))$
and $A**B = (\chi\ i\ j. \text{ inner } (A\$i) (\text{transpose } B\$j))$
proof –
have $(\sum_{k \in UNIV}. A \$i\$k * B \$k\$j) = \text{ inner } (\text{row } i\ A) (\text{column } j\ B)$
for $i\ j$

unfolding *column-def row-def inner-vec-def inner-real-def*
using *UNIV-I sum.cong vec-lambda-inverse* **by** *force*
thus $c1: A**B = (\chi\ i\ j.\ inner\ (row\ i\ A)\ (column\ j\ B))$
by *(simp add: matrix-matrix-mult-def)*
show $A**B = (\chi\ i\ j.\ inner\ (A\$i)\ (transpose\ B\$j))$
proof –
have $(\chi\ i\ j.\ A\ \$\ i \cdot transpose\ B\ \$\ j) = (\chi\ i\ j.\ row\ i\ A \cdot column\ j\ B)$
by *(simp add: row-def column-def transpose-def)*
then show *?thesis*
using *c1* **by** *metis*
qed
qed

lemma *matrix-inner-sum*:
shows $x \cdot y = (\sum\ i \in UNIV.\ \sum\ j \in UNIV.\ (x\$i\$j) \cdot (y\$i\$j))$
and $x \cdot y = (\sum\ (i,j) \in UNIV.\ (x\$i\$j) \cdot (y\$i\$j))$
apply *(simp add: inner-vec-def)+*
by *(simp add: sum.cartesian-product)*

lemma *matrix-norm-sum-sqrs*:
shows $norm\ x = sqrt(\sum\ i \in UNIV.\ \sum\ j \in UNIV.\ (norm\ (x\$i\$j))^2)$
and $norm\ x = sqrt(\sum\ (i,j) \in UNIV.\ (norm\ (x\$i\$j))^2)$
using *real-sqrt-abs real-sqrt-power*
by *(auto simp: norm-vec-def L2-set-def sum-nonneg sum.cartesian-product)*

lemma *norm-transpose*:
shows $norm\ x = norm\ (transpose\ x)$
proof –
have $(\sum\ (i,j) \in UNIV.\ (norm\ (x\$i\$j))^2) = (\sum\ (j,i) \in UNIV.\ (norm\ (x\$i\$j))^2)$
using *sum.swap[of $\lambda i j.\ (norm\ (x\$i\$j))^2$ UNIV UNIV]* **by** *(simp add: sum.cartesian-product)*
then show *?thesis*
unfolding *transpose-def matrix-norm-sum-sqrs(2)* **by** *simp*
qed

lemma *matrix-norm-inner*:
fixes $x::real^{n \times m}$
shows $norm\ x = sqrt(\sum\ (i,j) \in UNIV.\ (x\$i\$j) \cdot (x\$i\$j))$
using *matrix-inner-sum(2)[of x]* **by** *(simp add: norm-eq-sqrt-inner)*

lemma *matrix-norm-row*:
shows $norm\ x = sqrt(\sum\ i \in UNIV.\ (norm\ (row\ i\ x))^2)$
unfolding *norm-vec-def L2-set-def row-def* **by** *simp*

lemma *matrix-norm-column*:

shows $\text{norm } x = \text{sqrt}(\sum_{j \in \text{UNIV}}. (\text{norm } (\text{column } j \ x))^2)$
using *matrix-norm-row norm-transpose row-transpose*
by (*metis (lifting) Finite-Cartesian-Product.sum-cong-aux*)

lemma *mat-mul-indexed*: $(A**B)\$i\$j = (\sum_{k \in \text{UNIV}}. A \$ i \$ k * B \$ k \$ j)$

using *matrix-matrix-mult-def vec-lambda-beta*
by (*metis (no-types, lifting) Finite-Cartesian-Product.sum-cong-aux*)

lemma *norm-matrix-mult-ineq*:

fixes $A :: \text{real}^{\wedge l} \wedge n$
and $B :: \text{real}^{\wedge m} \wedge l$
shows $\text{norm } (A ** B) \leq \text{norm } A * \text{norm } B$

proof –

have $(A**B)\$i\$j = \text{row } i \ A \cdot \text{column } j \ B$ **for** $i \ j$
by (*simp add: matrix-matrix-mult-component-real(1)[of A B]*)
then have $\text{norm } (A**B) = \text{sqrt}(\sum_{(i,j) \in \text{UNIV}}. (\text{norm } (\text{row } i \ A \cdot \text{column } j \ B))^2)$
by (*simp add: matrix-norm-sum-sqrs(2)[of A**B]*)
then have $(\text{norm } (A**B))^2 = (\sum_{(i,j) \in \text{UNIV}}. (\text{norm } (\text{row } i \ A \cdot \text{column } j \ B))^2)$
by (*metis (no-types, lifting) norm-ge-zero real-sqrt-ge-0-iff real-sqrt-pow2*)
also have $(\sum_{(i,j) \in \text{UNIV}}. (\text{norm } (\text{row } i \ A \cdot \text{column } j \ B))^2)$
 $\leq (\sum_{(i,j) \in \text{UNIV}}. (\text{norm}(\text{row } i \ A) * \text{norm}(\text{column } j \ B))^2)$

proof –

obtain $f \ g$ **where** *defs*:

$f = (\lambda(i::'n, j::'m). (\text{row } i \ A \cdot \text{column } j \ B)^2)$
 $g = (\lambda(i::'n, j::'m). (\text{norm } (\text{row } i \ A) * \text{norm } (\text{column } j \ B))^2)$

by *simp*

then have $f \ (i, j) \leq g \ (i, j)$ **for** $i::'n$ **and** $j::'m$

by (*simp add: Cauchy-Schwarz-ineq power2-norm-eq-inner power-mult-distrib*)

hence $(\sum_{(i,j) \in \text{UNIV}}. f \ (i, j)) \leq (\sum_{(i,j) \in \text{UNIV}}. g \ (i, j))$

using *sum-mono[of UNIV f g]* **by** *fastforce*

thus *?thesis*

by (*simp add: defs*)

qed

also have $(\sum_{(i,j) \in \text{UNIV}}. (\text{norm}(\text{row } i \ A) * \text{norm}(\text{column } j \ B))^2) = (\text{norm } A * \text{norm } B)^2$

proof –

let $?f = \lambda i. (\text{norm } (\text{row } i \ A))^2$

let $?g = \lambda j. (\text{norm } (\text{column } j \ B))^2$

have $(\sum_{(i,j) \in \text{UNIV}}. (\text{norm}(\text{row } i \ A) * \text{norm}(\text{column } j \ B))^2)$
 $= (\sum_{(i,j) \in \text{UNIV}}. (\text{norm}(\text{row } i \ A))^2 * (\text{norm}(\text{column } j \ B))^2)$

by (*simp add: power-mult-distrib*)

then have $1: (\sum_{(i,j) \in \text{UNIV}}. (\text{norm}(\text{row } i \ A) * \text{norm}(\text{column } j \ B))^2)$
 $= (\sum_{i \in \text{UNIV}}. (\text{norm}(\text{row } i \ A))^2) * (\sum_{j \in \text{UNIV}}. (\text{norm}(\text{column } j \ B))^2)$

by (*simp add: sum-product sum.cartesian-product*)

have $2: (\sum_{i \in \text{UNIV}}. (\text{norm}(\text{row } i \ A))^2) = (\text{norm } A)^2 (\sum_{j \in \text{UNIV}}. (\text{norm}(\text{column } j \ B))^2) = (\text{norm } B)^2$

```

    using matrix-norm-row matrix-norm-column abs-norm-cancel real-sqrt-abs
real-sqrt-eq-iff
    by (smt (verit, best) sum.cong)+
    show ?thesis
    using 1 2 by (metis power-mult-distrib)
qed
finally show ?thesis
    by simp
qed

```

```

lemma bounded-bilinear-matrix-mult: bounded-bilinear (**)
  :: reallm ⇒ realnl ⇒ realnm
  apply (rule bounded-bilinear.intro)
  apply (metis (no-types, lifting) matrix-eq matrix-vector-mul-assoc matrix-vector-mult-add-rdistrib)
  apply (simp add: matrix-add-ldistrib matrix-scalar-ac scalar-matrix-assoc)+
  by (intro exI[of - 1], simp add: norm-matrix-mult-ineq)

```

```

lemma smooth-on-matrix-mult:
  fixes f::'a::real-normed-vector ⇒ (realnm)
  assumes k-smooth-on S f k-smooth-on S g open S
  shows k-smooth-on S (λx. f x ** g x)
  by (rule bounded-bilinear.smooth-on[OF bounded-bilinear-matrix-mult assms])

```

4.2 Smoothness of \prod and \det

```

lemma higher-differentiable-on-prod:
  fixes f::- ⇒ - ⇒ 'c::{real-normed-algebra, comm-monoid-mult}
  assumes  $\bigwedge i. i \in F \implies \text{finite } F \implies \text{higher-differentiable-on } S (f i) n \text{ open } S$ 
  shows higher-differentiable-on S (λx.  $\prod_{i \in F}. f i x$ ) n
  using assms apply (induction F rule: infinite-finite-induct)
  by (simp add: higher-differentiable-on-const higher-differentiable-on-mult)+

```

```

lemma smooth-on-prod:
  fixes f::- ⇒ - ⇒ 'c::{real-normed-algebra, comm-monoid-mult}
  assumes ( $\bigwedge i. i \in F \implies \text{finite } F \implies k\text{-smooth-on } S (f i)$ ) open S
  shows k-smooth-on S (λx.  $\prod_{i \in F}. f i x$ )
  using higher-differentiable-on-prod by (metis assms smooth-on-def)

```

```

lemma smooth-on-det:
  fixes s::('a::real-normed-field,'n::finite)square-matrix set
  assumes open s
  shows k-smooth-on s det
proof (unfold det-def, intro smooth-on-sum)
  fix p
  assume p ∈ {p. p permutes (UNIV::'n set)}
  show k-smooth-on s (λA. of-int (sign p) * ( $\prod_{i \in UNIV}. A \$ i \$ p i$ ))
  proof (intro smooth-on-mult)
    show k-smooth-on s (λx. of-int (sign p))

```

```

    by (simp add: smooth-on-const)
  show k-smooth-on s (λx. ∏ i ∈ UNIV. x $ i $ p i) open s
    apply (intro smooth-on-prod)
    apply (simp add: bounded-linear.smooth-on bounded-linear-mat-ijth)
    by (rule assms)+
  qed
qed (rule assms)

```

4.3 Smoothness of matrix inversion

```

lemma invertible-mat-1: invertible (mat 1)
  by (simp add: invertible-def)

```

```

lemma continuous-on-vec:
  assumes ∧i. continuous-on S (λx. f x $ i)
  shows continuous-on S f
  using assms unfolding continuous-on-def by (simp add: vec-tendstoI)

```

```

lemma frechet-derivative-eucl:
  fixes f::'a::euclidean-space ⇒ 'b::real-normed-vector
  assumes f differentiable at x
  shows frechet-derivative f (at x) =
    (λv. ∑ i ∈ Basis. (v · i) *R frechet-derivative f (at x) i)
proof -
  have 1: id differentiable at x f differentiable at (id x)
    by (simp add: frechet-derivative-works, simp add: assms)
  show ?thesis using frechet-derivative-compose-eucl[OF 1] frechet-derivative-id[of
x]
    by (auto, metis comp-id fun.map-ident)
qed

```

TODO! This should maybe be changed in *Finite-Cartesian-Product.norm-le-l1-cart*. That result only works for real^n , this one should work for all $'a::\text{real-normed-vector}^n$.

```

lemma norm-le-l1-cart': norm x ≤ sum(λi. norm (x $ i)) UNIV
  by (simp add: norm-vec-def L2-set-le-sum)

```

```

lemma bounded-linear-vec-nth-fun:
  fixes f::'a::real-normed-vector ⇒ 'b::real-normed-vectorm
  assumes ∧i. bounded-linear (λx. (f x) $ i)
  shows bounded-linear f
proof
  fix x y and r::real
  interpret fi: bounded-linear λx. (f x) $ i for i by fact
  show f (r *R x) = r *R f x
    using fi.scale by (simp add: vec-eq-iff)
  show f (x + y) = f x + f y
    using fi.add by (simp add: vec-eq-iff)
  obtain F where 0 < F i and norm-f: ∧x. norm ((f x) $ i) ≤ norm x * F i for i
    using fi.pos-bounded by metis

```

```

have  $\forall x. \text{norm } (f x) \leq \text{norm } x * (\sum_{i \in \text{UNIV}} F i)$ 
proof (rule allI)
  fix x
  have  $\text{norm } (f x) \leq (\sum_{i \in \text{UNIV}} \text{norm } (f x \$ i))$ 
    by (rule norm-le-l1-cart'[of f x for x])
  also have  $\dots \leq (\sum_{i \in \text{UNIV}} \text{norm } x * F i)$ 
    using norm-f[of x i for i] by (simp add: sum-mono)
  also have  $\dots \leq \text{norm } x * (\sum_{i \in \text{UNIV}} F i)$ 
    by (simp add: sum-distrib-left)
  finally show  $\text{norm } (f x) \leq \text{norm } x * (\sum_{i \in \text{UNIV}} F i)$  .
qed
thus  $\exists K. \forall x. \text{norm } (f x) \leq \text{norm } x * K$  by blast
qed

```

```

lemma has-derivative-vec-lambda [derivative-intros]:
  fixes  $f::'a::\text{real-normed-vector} \Rightarrow 'b::\text{real-normed-vector}^m$ 
  assumes  $\bigwedge i. ((\lambda x. (f x)\$i) \text{ has-derivative } (\lambda x. (f' x)\$i))$  (at x within s)
  shows (f has-derivative f') (at x within s)
proof (intro has-derivativeI-sandwich[of 1])
  show bounded-linear f'
    using assms by (intro bounded-linear-vec-nth-fun has-derivative-bounded-linear)

  let ?Ri =  $\lambda i y. (f y)\$i - (f x)\$i - (f' (y-x))\$i$ 
  let ?R =  $\lambda y. f y - f x - f' (y-x)$ 

  show  $((\lambda y. (\sum_{i \in \text{UNIV}} \text{norm } (?Ri i y) / \text{norm } (y-x))) \longrightarrow 0)$  (at x within s)
    using assms apply (intro tendsto-null-sum) by (auto simp: has-derivative-iff-norm)

  fix y assume  $y \neq x$ 
  show  $\text{norm } (?R y) / \text{norm } (y-x) \leq (\sum_{i \in \text{UNIV}} \text{norm } (?Ri i y) / \text{norm } (y-x))$ 
    unfolding sum-divide-distrib[symmetric]
    apply (rule divide-right-mono) prefer 2 apply simp
    using norm-le-l1-cart' by (smt (verit, ccfv-SIG) real-norm-def sum-mono vector-minus-component)
qed (simp)

```

```

lemma has-derivative-vec-lambda-2:
  fixes  $f::'a::\text{real-normed-vector} \Rightarrow 'b::\text{real-normed-vector}^m$ 
  assumes  $\bigwedge i. ((\lambda x. (f x)\$i) \text{ has-derivative } (f' i))$  (at x within s)
  shows (f has-derivative  $(\lambda x. \chi i. f' i x)$ ) (at x within s)
  apply (intro has-derivative-vec-lambda[of f  $\lambda x. \chi i. f' i x$  s])
  using assms by auto

```

```

lemma differentiable-componentwise:
  fixes  $f::'a::\text{real-normed-vector} \Rightarrow 'b::\text{real-normed-vector}^m$ 
  assumes  $\bigwedge i. (\lambda x. f x \$ i) \text{ differentiable}$  (at x within s)
  shows f differentiable (at x within s)
proof (unfold differentiable-def, intro exI)

```

```

let ?f' = λi. SOME f'. ((λx. f x $ i) has-derivative f') (at x within s)
have 1: ∧i. ((λx. (f x) $ i) has-derivative (?f' i)) (at x within s)
  by (metis assms differentiable-def some-eq-imp)
show (f has-derivative (λx. χ i. ?f' i x)) (at x within s)
  by (rule has-derivative-vec-lambda-2[OF 1])
qed

lemma frechet-derivative-vec:
  fixes f::'a::real-normed-vector ⇒ 'b::real-normed-vector^'m
  assumes ∧i. (λx. f x $ i) differentiable (at x)
  shows frechet-derivative f (at x) = (λv. χ i. (frechet-derivative (λx. f x $ i) (at x) v))
  apply (rule frechet-derivative-at')
  apply (intro has-derivative-vec-lambda)
  by (auto intro: derivative-eq-intros frechet-derivative-worksI[OF assms])

lemma higher-differentiable-on-vec:
  fixes f::'a::real-normed-vector ⇒ 'b::real-normed-vector^'m
  assumes ∧i. higher-differentiable-on S (λx. (f x) $ i) n
  and open S
  shows higher-differentiable-on S f n
  using assms
proof (induction n arbitrary: f)
  case 0
  then show ?case
    using continuous-on-vec by (metis higher-differentiable-on.simps(1))
next
  case (Suc n)
  have f: ∧x i. x ∈ S ⇒ (λx. f x $ i) differentiable (at x)
  and hf: ∧i. higher-differentiable-on S (λx. frechet-derivative (λy. f y $ i) (at x) v) n
  for v using Suc.prem higher-differentiable-on.simps(2) by blast+
  have f': higher-differentiable-on S f n
  using Suc(1,2) assms(2) higher-differentiable-on-SucD by blast
  have 1: ∀x∈S. f differentiable (at x)
  using f differentiable-componentwise[of f - UNIV] by simp
  have 2: ∀v. higher-differentiable-on S (λx. frechet-derivative f (at x) v) n
proof (intro allI)
  fix v
  let ?f' = λx. frechet-derivative f (at x) v
  let ?f'_i = λx. χ i. frechet-derivative (λy. f y $ i) (at x) v
  { fix x assume x ∈ S
    hence ?f' x = (χ i. frechet-derivative (λy. f y $ i) (at x) v)
    using frechet-derivative-vec[OF f] by simp }
  then have higher-differentiable-on S ?f' n = higher-differentiable-on S ?f'_i n
  using higher-differentiable-on-cong[of S S ?f' ?f'_i n] assms(2)
  by simp
  then show higher-differentiable-on S ?f' n
  using hf Suc.IH assms(2) by auto

```

```

qed
show ?case
  by (simp add: 1 2 higher-differentiable-on.simps(2))
qed

```

```

lemma smooth-on-vec:
  fixes f::'a::real-normed-vector  $\Rightarrow$  'b::real-normed-vector^m
  assumes  $\bigwedge i. k\text{-smooth-on } S (\lambda x. (f x) \$ i)$  open S
  shows k-smooth-on S f
proof (unfold smooth-on-def, intro allI impI)
  fix n assume asm: enat n  $\leq$  k
  show higher-differentiable-on S f n
  apply (intro higher-differentiable-on-vec)
  using assms asm unfolding smooth-on-def by simp+
qed

```

```

lemma smooth-on-mat:
  fixes f::('a::real-normed-vector)  $\Rightarrow$  ('b::real-normed-vector^k)^l
  assumes  $\bigwedge i j. k\text{-smooth-on } S (\lambda x. (f x)\$i\$j)$  open S
  shows k-smooth-on S f
  by (simp add: smooth-on-vec assms)

```

This type constraint is annoying. The *euclidean-space* is inherited from *higher-differentiable-on-compose*, where it is marked as: ‘TODO: can we get around this restriction’. Notice this type constraint is exactly *real-normed-eucl* as defined in *Classical-Groups*.

```

lemma smooth-on-matrix-inv-component:
  fixes S::('a::{euclidean-space,real-normed-field})^n^n set
  assumes  $\forall A \in S. \text{invertible } A$  open S
  shows k-smooth-on S ( $\lambda A. (\text{matrix-inv } A)\$i\$j$ )
  using matrix-inverse-component smooth-on-mat smooth-on-det smooth-on-compose
  smooth-on-divide smooth-on-cong
proof -
  have smooth-on-div-det: k-smooth-on S ( $\lambda x. f x / (\det x)$ ) if smooth-on S f for
  f
  apply (intro smooth-on-divide[of k S f det])
  using that smooth-on-det[OF assms(2)] assms by (auto simp: smooth-on-def
  invertible-det-nz)

```

```

  let ?inv-comp' =  $\lambda A::'a^{\wedge}n^{\wedge}n. \chi k l. \text{if } k = j \wedge l = i \text{ then } 1 \text{ else if } k = j \vee l = i \text{ then } 0 \text{ else } A \$ k \$ l$ 

```

```

  let ?inv-comp =  $\lambda A::'a^{\wedge}n^{\wedge}n. \det (?inv-comp' A) / \det A$ 

```

```

  have matrix-inv-cong:  $\bigwedge A. A \in S \implies (\text{matrix-inv } A)\$i\$j = ?inv-comp A$ 
  using matrix-inverse-component assms by blast

```

```

  have smooth-on-component: smooth-on S ?inv-comp'
proof (intro smooth-on-mat[of  $\infty$  S ?inv-comp'])
  fix n m

```



```

consider  $n=j \wedge m=i \mid n=j \wedge m \neq i \mid n \neq j \wedge m=i \mid n \neq j \wedge m \neq i$  by linarith
hence smooth-on S ( $\lambda x.$  if  $n = j \wedge m = i$  then 1 else if  $n = j \vee m = i$  then 0
else  $x \$ n \$ m$ )
  apply cases by (simp add: smooth-on-const smooth-on-ijth-of-mat) +
  thus smooth-on S ( $\lambda x.$  ( $\chi$   $k$   $l.$  if  $k = j \wedge l = i$  then 1 else if  $k = j \vee l = i$  then
0 else  $x \$ k \$ l$ )  $\$ n \$ m$ )
    by simp
  qed (fact)

thus k-smooth-on S ( $\lambda A.$  (matrix-inv A) $\$i\$j$ )
  apply (intro smooth-on-cong[OF - assms(2) matrix-inv-cong])
  apply (intro smooth-on-div-det[of  $\lambda A.$  det (?inv-comp' A)])
  using smooth-on-compose[of  $\infty$  UNIV det S ?inv-comp'] smooth-on-det[OF
open-UNIV]
  using assms(2) smooth-on-cong by fastforce
qed

```

```

lemma fin-sum-over-delta:
  fixes  $f::'n::\text{finite} \Rightarrow 'a::\text{semiring-1}$ 
  shows ( $\sum (i::'n::\text{finite}) \in \text{UNIV}.$  (if  $i=j$  then 1 else 0) *  $f$   $i$ ) =  $f$   $j$ 
proof -
  have ( $\sum i \in \text{UNIV}.$  (if  $i = j$  then 1 else 0) *  $f$   $i$ ) = ( $\sum i \in \text{UNIV}.$  (if  $i=j$  then  $f$   $j$ 
else 0))
    by (simp add: mult-delta-left)
  also have ( $\sum i \in \text{UNIV}.$  (if  $i=j$  then  $f$   $j$  else 0)) =  $f$   $j$ 
    using sum.delta by auto
  then show ?thesis
    by (simp add: calculation)
qed

```

```

lemma matrix-is-linear-map:
  fixes  $A::('a::\{\text{real-algebra-1, comm-semiring-1}\})^m \wedge n$  — again, real-based entries
only...
  shows linear (( $*$ ) $v$ )  $A$   $\wedge$  matrix (( $*$ ) $v$ )  $A$  =  $A$ 
proof (rule conjI)
  let  $?f = \lambda v.$  ( $A$  *  $v$ )
  show linear ?f
    using matrix-vector-mul-linear by simp
  {
    fix  $i::'n$  and  $j::'m$ 
    let  $?v = \chi$   $j'.$  if  $j' = j$  then 1 else 0
    have  $?v \$ k =$  (if  $k=j$  then 1 else 0) for  $k$ 
      by simp
    then have  $A * v$   $?v =$  transpose A  $\$ j$ 
      using matrix-vector-column[where  $A=A$  and  $x=?v$ ] fin-sum-over-delta
    by (smt (verit, best) mult.commute mult.right-neutral mult-zero-right sum.cong
vector-smult-lid vector-smult-lzero)
  }

```

```

then have (A *v (χ j'. if j' = j then 1 else 0))$i = A$i$j
using matrix-vector-column[where x=?v] transpose-def vec-lambda-beta
by (smt (z3))
}
then show matrix ?f = A
unfolding matrix-def axis-def by auto
qed

```

```

lemma smooth-on-matrix-inv:
assumes ∀ A. A ∈ S → invertible A open S
shows k-smooth-on S (matrix-inv::'a::{euclidean-space,real-normed-field} ^n ^n
⇒ 'a ^n ^n)
apply (intro smooth-on-mat[of k S])
apply (intro smooth-on-matrix-inv-component[of S])
by (auto simp add: assms)+

```

end

5 Smooth vector fields

theory Smooth-Vector-Fields

imports

More-Manifolds

begin

Type synonyms for use later: these already follow our later split between defining “charts” for the tangent bundle as a product, and talking about vector fields as maps $p \mapsto v \in T_p M$ as well as sections of the tangent bundle $M \rightarrow TM$.

type-synonym 'a tangent-bundle = 'a × (('a ⇒ real) ⇒ real)

type-synonym 'a vector-field = 'a ⇒ (('a ⇒ real) ⇒ real)

5.1 (Smooth) vector fields on an (entire) manifold.

Since we only get an isomorphism between tangent vectors and directional derivatives in the smooth case of $k = \infty$, we create a locale for infinitely smooth manifolds.

locale smooth-manifold = c-manifold charts ∞ **for** charts

context c-manifold **begin**

5.1.1 Charts for the tangent bundle

definition in-TM :: 'a ⇒ (('a ⇒ real) ⇒ real) ⇒ bool

where in-TM p v ≡ p ∈ carrier ∧ v ∈ tangent-space p

abbreviation $TM \equiv \{(p,v). \text{in-TM } p \ v\}$

lemma in-TM-E [*elim*]:
assumes $\text{in-TM } p \ v$
shows $v \in \text{tangent-space } p \ p \in \text{carrier}$
using *assms* **unfolding** in-TM-def **by** *auto*

lemma $TM\text{-PairE}$ [*elim*]:
assumes $(p,v) \in TM$
shows $v \in \text{tangent-space } p \ p \in \text{carrier}$
using *assms* **unfolding** in-TM-def **by** *auto*

lemma $TM\text{-E}$ [*elim*]:
assumes $x \in TM$
shows $\text{snd } x \in \text{tangent-space } (\text{fst } x) \ \text{fst } x \in \text{carrier}$
using *assms* **by** *auto*

We can construct a chart for *tangent-space* p given a chart around p . Notice the appearance of *charts* in the definition, which specifies that we're charting the set *tangent-space* p , not *c-manifold.tangent-space* (*charts-submanifold* c) ∞ p .

definition $\text{apply-chart-TM} :: ('a,'b)\text{chart} \Rightarrow 'a \text{ tangent-bundle} \Rightarrow 'b \times 'b$
where $\text{apply-chart-TM } c \equiv \lambda(p,v). (c \ p \ , \ c\text{-manifold-point.tangent-chart-fun } c \ p \ v)$

definition $\text{inv-chart-TM} :: ('a,'b)\text{chart} \Rightarrow ('b \times 'b) \Rightarrow 'a \times (('a \Rightarrow \text{real}) \Rightarrow \text{real})$
where $\text{inv-chart-TM } c \equiv \lambda((p::'b),(v::'b)). (\text{inv-chart } c \ p \ , \ c\text{-manifold-point.coordinate-vector } c \ (\text{inv-chart } c \ p) \ v)$

definition $\text{domain-TM} :: ('a,'b) \text{chart} \Rightarrow ('a \times (('a \Rightarrow \text{real}) \Rightarrow \text{real})) \text{ set}$
where $\text{domain-TM } c \equiv \{(p, v). p \in \text{domain } c \wedge v \in \text{tangent-space } p\}$

definition $\text{codomain-TM} :: ('a,'b) \text{chart} \Rightarrow ('b \times 'b) \text{ set}$
where $\text{codomain-TM } c \equiv \{(p, v). p \in \text{codomain } c\}$

definition $\text{restrict-chart-TM } S \ c \equiv \text{apply-chart-TM } (\text{restrict-chart } S \ c)$

definition $\text{restrict-domain-TM } S \ c \equiv \text{domain-TM } (\text{restrict-chart } S \ c)$

definition $\text{restrict-codomain-TM } S \ c \equiv \text{codomain-TM } (\text{restrict-chart } S \ c)$

definition $\text{restrict-inv-chart-TM } S \ c \equiv \text{inv-chart-TM } (\text{restrict-chart } S \ c)$

5.1.2 Proofs about apply-chart-TM that mimic the properties of $('a, 'b) \text{chart}$.

lemma domain-TM :
assumes $c \in \text{atlas}$
shows $\text{domain-TM } c \subseteq TM$
unfolding domain-TM-def in-TM-def **using** *assms* **by** *auto*

lemma codomain-TM-alt : $\text{codomain-TM } c = \text{codomain } c \times (\text{UNIV} :: 'b \text{ set})$

```

unfolding codomain-TM-def by auto

lemma open-codomain-TM:
  assumes  $c \in \text{atlas}$ 
  shows open (codomain-TM  $c$ )
  using codomain-TM-alt open-Times[OF open-codomain open-UNIV] by auto

end

context smooth-manifold begin

lemma apply-chart-TM-inverse [simp]:
  assumes  $c: c \in \text{atlas}$ 
  shows  $\bigwedge p v. (p,v) \in \text{domain-TM } c \implies \text{inv-chart-TM } c (\text{apply-chart-TM } c (p,v)) = (p,v)$ 
  and  $\bigwedge x u. (x,u) \in \text{codomain-TM } c \implies \text{apply-chart-TM } c (\text{inv-chart-TM } c (x,u)) = (x,u)$ 
  proof –
    fix  $p v$  assume  $(p,v) \in \text{domain-TM } c$ 
    then have asm:  $c \in \text{atlas } p \in \text{domain } c v \in \text{tangent-space } p$ 
      using  $c$  by (auto simp add: domain-TM-def)
    interpret  $p$ : c-manifold-point charts  $\infty$   $c p$ 
      using c-manifold-point[OF asm(1,2)] by simp
    have  $v \in p.T_p M$  using asm(3) by simp
    from  $p.\text{coordinate-vector-inverse}(1)$ [OF - this] show  $\text{inv-chart-TM } c (\text{apply-chart-TM } c (p,v)) = (p,v)$ 
      by (simp add: inv-chart-TM-def apply-chart-TM-def p.tangent-chart-fun-def)
    next
      fix  $x u$  assume  $(x,u) \in \text{codomain-TM } c$ 
      then have asm:  $c \in \text{atlas } x \in \text{codomain } c$ 
        using  $c$  by (auto simp add: codomain-TM-def)
      interpret  $x$ : c-manifold-point charts  $\infty$   $c \text{inv-chart } c x$ 
        using c-manifold-point[OF asm(1)] by (simp add: asm(2))
      from  $x.\text{coordinate-vector-inverse}(2)$  show  $\text{apply-chart-TM } c (\text{inv-chart-TM } c (x,u)) = (x,u)$ 
        by (simp add: inv-chart-TM-def apply-chart-TM-def x.tangent-chart-fun-def asm(2))
    qed

lemma image-domain-TM-eq:
  assumes  $c \in \text{atlas}$ 
  shows  $\text{apply-chart-TM } c \text{ ` domain-TM } c = \text{codomain-TM } c$ 
  proof –
    { fix  $x :: 'b \times 'b$  assume  $x: x \in \text{codomain } c \times \text{UNIV}$ 
      obtain  $y_1 y_2$  where  $y_1 = \text{inv-chart } c (\text{fst } x) y_2 = \text{c-manifold-point.coordinate-vector charts } \infty$   $c y_1 (\text{snd } x)$ 
      by simp
      have  $y_1 \in \text{domain } c$  using  $y(1) x$  by auto
    }

```

then interpret y_1 : *c-manifold-point charts* ∞ c y_1
by (*simp add: assms(1) c-manifold-point*)
have $y_2 \in$ *tangent-space* y_1
using $y(2)$ x *assms* y_1 .*coordinate-vector-surj* **by** *blast*
then have $(y_1, y_2) \in \{(p, v). p \in$ *domain* $c \wedge v \in$ *tangent-space* $p\}$
using $\langle y_1 \in$ *domain* $c \rangle$ **by** *simp*
moreover have *fst* $x = c$ y_1 *snd* $x = c$ -*manifold-point.tangent-chart-fun charts*
 ∞ c y_1 y_2
using y x *assms* y_1 .*tangent-chart-fun-inverse(2)* **by** *auto*
ultimately have $x \in (\lambda(p, v). (c$ p, c -*manifold-point.tangent-chart-fun charts*
 ∞ c p $v)) \{(p, v). p \in$ *domain* $c \wedge v \in$ *tangent-space* $p\}$
by (*metis (no-types, lifting) pair-imageI prod.collapse*) }
thus *?thesis* **by** (*auto simp: apply-chart-TM-def domain-TM-def codomain-TM-alt*)
qed

lemma *inv-image-codomain-TM-eq*:
assumes $c \in$ *atlas*
shows *inv-chart-TM* c ‘ *codomain-TM* $c =$ *domain-TM* c
apply (*subst image-domain-TM-eq[OF assms, symmetric]*)
using *apply-chart-TM-inverse(1)[OF assms]* **by** *force*

lemma (*in c-manifold*) *restrict-domain-TM-intersection*:
shows *restrict-domain-TM* (*domain* $c1 \cap$ *domain* $c2$) $c1 =$ *domain-TM* $c1 \cap$
domain-TM $c2$
unfolding *restrict-domain-TM-def* **by** (*auto simp: domain-TM-def open-Int*)

lemma (*in c-manifold*) *restrict-domain-TM-intersection'*:
shows *restrict-domain-TM* (*domain* $c1 \cap$ *domain* $c2$) $c2 =$ *domain-TM* $c1 \cap$
domain-TM $c2$
unfolding *restrict-domain-TM-def* **by** (*auto simp: domain-TM-def open-Int*)

lemma (*in c-manifold*) *restrict-domain-TM*:
assumes *open* S $S \subseteq$ *domain* c
shows *restrict-domain-TM* S $c = \{(p, v). p \in$ $S \wedge v \in$ *tangent-space* $p\}$
unfolding *restrict-domain-TM-def domain-TM-def* **using** *domain-restrict-chart*
assms **by** *auto*

lemma *image-restrict-domain-TM-eq*:
assumes $c \in$ *atlas*
shows *restrict-chart-TM* S c ‘ *restrict-domain-TM* S $c =$ *restrict-codomain-TM*
 S c
unfolding *restrict-chart-TM-def restrict-domain-TM-def restrict-codomain-TM-def*
using *image-domain-TM-eq assms restrict-chart-in-atlas* **by** *blast*

lemma *inv-image-restrict-codomain-TM-eq*:
assumes $c \in \text{atlas}$
shows $\text{restrict-inv-chart-TM } S \ c \ ' \ \text{restrict-codomain-TM } S \ c = \text{restrict-domain-TM } S \ c$
by (*metis (no-types, lifting) inv-image-codomain-TM-eq assms restrict-chart-in-atlas restrict-codomain-TM-def restrict-domain-TM-def restrict-inv-chart-TM-def*)

lemma *codomain-restrict-chart-TM[simp]*:
assumes $c \in \text{atlas}$ *open* S
shows $\text{restrict-codomain-TM } S \ c = \text{codomain-TM } c \cap \text{inv-chart-TM } c \ - \{ (p, v). p \in S \wedge v \in \text{tangent-space } p \}$
proof –
{
 fix $a \ b \ p \ v$
 assume $asm: a \in \text{codomain } c \ \text{inv-chart-TM } c \ (a, b) = (p, v)$
 interpret $p: c\text{-manifold-point charts } \infty \ c \ \text{inv-chart } c \ a$
 using $asm(1) \ \text{assms } c\text{-manifold-point}[OF \ \text{assms}(1), \ \text{of inv-chart } c \ a \ \text{for } a]$ **by** *blast*
 have $p.\text{coordinate-vector } b \in \text{tangent-space } (\text{inv-chart } c \ a)$
 using $\text{bij-betweE}[OF \ p.\text{coordinate-vector-bij}]$ **by** *simp*
 then have $\text{inv-chart } c \ a \in S \implies v \in \text{tangent-space } p$
 and $\text{inv-chart } c \ a \in S \implies p \in S$
 and $\llbracket p \in S; v \in \text{tangent-space } p \rrbracket \implies \text{inv-chart } c \ a \in S$
 subgoal using $\text{inv-chart-TM-def inv-image-codomain-TM-eq}[OF \ \text{assms}(1)]$
asm **by** *auto*
 subgoal using $asm(2)$ **by** (*auto simp add: assms(2) inv-chart-TM-def*)
 subgoal using $asm(2) \ c\text{-manifold.inv-chart-TM-def}[OF \ c\text{-manifold-axioms}]$
by *simp*
 done
}
thus *?thesis* **by** (*auto simp add: restrict-codomain-TM-def codomain-TM-def assms(2)*)
qed

lemma (*in* $c\text{-manifold}$) *image-subset-TM-eq [simp]*:
assumes $S \subseteq \text{domain-TM } c$
shows $\text{apply-chart-TM } c \ ' \ S \subseteq \text{codomain-TM } c$
using assms **unfolding** $\text{apply-chart-TM-def codomain-TM-def domain-TM-def}$
by *auto*

lemma (*in* $c\text{-manifold}$) *image-subset-restrict-TM-eq [simp]*:
assumes $T \subseteq \text{restrict-domain-TM } S \ c$
shows $\text{restrict-chart-TM } S \ c \ ' \ T \subseteq \text{restrict-codomain-TM } S \ c$
using assms **unfolding** $\text{restrict-chart-TM-def restrict-codomain-TM-def restrict-domain-TM-def}$
by *auto*

lemma *restrict-chart-domain-Int*:

```

assumes  $c1 \in \text{atlas}$ 
shows  $\text{apply-chart-TM } c1 \text{ ' (domain-TM } c1 \cap \text{domain-TM } c2) = \text{restrict-chart-TM}$ 
 $(\text{domain } c1 \cap \text{domain } c2) \text{ } c1 \text{ ' (restrict-domain-TM (domain } c1 \cap \text{domain } c2) \text{ } c1)$ 
 $(\text{is } \langle ?\text{TM-dom-Int} = ?\text{restr-TM-dom} \rangle)$ 
proof (intro subset-antisym)
  have  $\text{dom-eq: domain (restrict-chart (domain } c1 \cap \text{domain } c2) \text{ } c1) = \text{domain } c1$ 
 $\cap \text{domain } c2$ 
    using  $\text{domain-restrict-chart}[OF \text{ open-domain}]$  by (metis inf.left-idem)

  { fix  $x$  assume  $x \in (\text{domain-TM } c1 \cap \text{domain-TM } c2)$ 
    then obtain  $p \ v$  where  $x: x = (p,v) \ p \in \text{domain } c1 \ p \in \text{domain } c2 \ v \in$ 
 $\text{tangent-space } p$ 
      unfolding  $\text{domain-TM-def}$  by blast
      interpret  $p1: c\text{-manifold-point charts } \infty \ c1 \ p$  using  $c\text{-manifold-point}[OF$ 
 $\text{assms}(1) \ x(2)]$  by simp
      interpret  $p2: c\text{-manifold-point charts } \infty \ \text{restrict-chart (domain } c1 \cap \text{domain}$ 
 $c2) \ c1 \ p$ 
      using  $c\text{-manifold-point}[OF \text{assms}(1) \ x(2)]$   $\text{restrict-chart-in-atlas}[OF \text{assms}(1)]$ 
 $\text{domain-restrict-chart}[OF \text{ open-domain}]$ 
      by (metis IntI c-manifold-point p1.p x(3))

    have  $[simp]: p2.\text{sub-}\psi.\text{sub.restrict-codomain-TM (domain } c1 \cap \text{domain } c2) \ c1$ 
 $=$ 
       $\{(p, v). p \in \text{codomain (restrict-chart (domain } c1 \cap \text{domain } c2) \ c1)\}$ 
    unfolding  $p2.\text{sub-}\psi.\text{sub.restrict-codomain-TM-def}$   $p2.\text{sub-}\psi.\text{sub.codomain-TM-def}$ 
by simp

    have  $\text{apply-chart-TM } c1 \ x \in ?\text{restr-TM-dom}$ 
      apply (simp add: x(1) image-restrict-domain-TM-eq[OF assms(1)])
      unfolding  $\text{apply-chart-TM-def}$  using  $p2.\psi p\text{-in}$  by (auto simp: p1.euclidean-coordinates-eq-iff)
    }
  thus  $?\text{TM-dom-Int} \subseteq ?\text{restr-TM-dom}$  by auto

  { fix  $x$  assume  $x \in \text{restrict-domain-TM (domain } c1 \cap \text{domain } c2) \ c1$ 
    then obtain  $p \ v$  where  $x: x = (p,v) \ p \in \text{domain } c1 \ p \in \text{domain } c2 \ v \in$ 
 $\text{tangent-space } p$ 
      unfolding  $\text{restrict-domain-TM-def}$   $\text{domain-TM-def}$  by (auto simp: dom-eq)
      interpret  $p1: c\text{-manifold-point charts } \infty \ c1 \ p$  using  $c\text{-manifold-point}[OF$ 
 $\text{assms}(1) \ x(2)]$  by simp
      interpret  $p2: c\text{-manifold-point charts } \infty \ \text{restrict-chart (domain } c1 \cap \text{domain}$ 
 $c2) \ c1 \ p$ 
      using  $c\text{-manifold-point}[OF \text{assms}(1) \ x(2)]$   $\text{restrict-chart-in-atlas}[OF \text{assms}(1)]$ 
 $\text{domain-restrict-chart}[OF \text{ open-domain}]$ 
      by (metis IntI c-manifold-point p1.p x(3))
    have  $\text{restrict-chart-TM (domain } c1 \cap \text{domain } c2) \ c1 \ x \in ?\text{TM-dom-Int}$ 
proof –
      have  $\text{apply-chart } c1 \ p \in \text{apply-chart } c1 \text{ ' (domain } c1 \cap \text{domain } c2)$ 
        using  $p1.p \ x(3)$  by blast
      moreover have  $p2.\text{tangent-chart-fun } v \in c\text{-manifold-point.tangent-chart-fun}$ 

```

```

charts ∞ c1 p ‘ {v. v∈tangent-space p}
  using p1.coordinate-vector-surj p1.tangent-chart-fun-inverse(2) by fastforce
  ultimately show ?thesis
  apply (simp add: apply-chart-TM-def)
  apply (simp add: x(1) restrict-chart-TM-def)
  apply (simp add: apply-chart-TM-def apply-chart-restrict-chart[of domain
c1 ∩ domain c2 c1])
  unfolding domain-TM-def by force
  qed }
thus ?restr-TM-dom ⊆ ?TM-dom-Int by blast
qed

```

```

lemma open-intersection-TM:
  assumes c1 ∈ atlas
  shows open (apply-chart-TM c1 ‘ (domain-TM c1 ∩ domain-TM c2))
  using restrict-chart-domain-Int image-restrict-domain-TM-eq restrict-chart-in-atlas
  assms
  by (auto simp: restrict-codomain-TM-def open-codomain-TM)

```

```

lemma apply-restrict-chart-TM:
  assumes c: c ∈ atlas and S: open S S ⊆ domain c x ∈ restrict-domain-TM S c
  shows apply-chart-TM c x = restrict-chart-TM S c x
proof –
  { fix p v assume x: x = (p,v) p ∈ S v ∈ tangent-space p
    interpret p1: c-manifold-point charts ∞ c p
      using c-manifold-point[OF c] x(2) S(2) by blast
    interpret p2: c-manifold-point charts ∞ restrict-chart S c p
      apply (rule c-manifold.c-manifold-point, unfold-locales)
      using S(1) x(2) by (auto simp add: restrict-chart-in-atlas)
    have TpM-eq: p2.TpM = tangent-space p by simp
    have p1.tangent-chart-fun v = p2.tangent-chart-fun v
      unfolding p1.tangent-chart-fun-def p2.tangent-chart-fun-def
      using p1.component-function-restrict-chart[OF x(2) S(1)] TpM-eq x(3) by
  simp }
  thus ?thesis
    using S(3) restrict-domain-TM[OF S(1,2)] unfolding restrict-chart-TM-def
  apply-chart-TM-def by auto
qed

```

```

lemma inverse-restrict-chart-TM:
  assumes c: c ∈ atlas and S: open S S ⊆ domain c x ∈ restrict-codomain-TM S
  c
  shows inv-chart-TM c x = restrict-inv-chart-TM S c x
proof –
  { fix p v assume x: x = (p,v) p ∈ c‘S
    interpret p1: c-manifold-point charts ∞ c inv-chart c p

```



```

    using c-manifold-point[OF c] x(2) S(2) by blast
  have pS: inv-chart c p ∈ S
    using restrict-chart-in-atlas x(2) S(2) image-domain-eq by auto
  interpret p2: c-manifold-point charts ∞ restrict-chart S c inv-chart c p
  apply (rule c-manifold.c-manifold-point, unfold-locales)
  using pS restrict-chart-in-atlas S(1) by auto
  have p1.coordinate-vector v = p2.coordinate-vector v
  using p1.coordinate-vector-restrict-chart[OF pS S(1)]
  using p1.coordinate-vector-def p2.coordinate-vector-def by presburger }
thus ?thesis
  using S(3) inv-chart-TM-def apply-chart-TM-def
  apply (simp add: codomain-restrict-chart-TM[OF c S(1)] restrict-inv-chart-TM-def)
  using apply-chart-TM-inverse(2)[OF c] surj-pair by (smt (verit) case-prod-conv
image-eqI)
qed

```

lemma (in *c-manifold-point*) *dκ-inv-directional-derivative-eq*:

```

  assumes k = ∞
  shows  $d\kappa^{-1}$  (directional-derivative k (ψ p) x) = restrict0 (diffeo-ψ.dest.diff-fun-space)
  (λf. frechet-derivative f (at (ψ p)) x)
proof –

```

```

  let ?is-ext = λf f'. f ∈ diffeo-ψ.dest.diff-fun-space ∧ f' ∈ manifold-eucl.dest.diff-fun-space
  ∧ f' ∈ diffeo-ψ.dest.diff-fun-space ∧
  (∃ N. ψ p ∈ N ∧ open N ∧ closure N ⊆ diffeo-ψ.dest.carrier ∧ (∀ x ∈ closure N.
  f' x = f x) ∧
  (∀ v ∈ Tψ p ψ U. v f = v f') ∧ (∀ v ∈ Tψ p ψ U. v f' = dκ v f') ∧ (∀ v ∈ Tψ p E. dκ-1
  v f = v f'))

```

```

  let ?extend = λf. SOME f'. ?is-ext f f'
  obtain extend where extend-def: extend ≡ ?extend by blast
  have extend: ?is-ext f (extend f) ?is-ext f (?extend f)
  if f ∈ diffeo-ψ.dest.diff-fun-space for f
proof –
  show ?is-ext f (?extend f)
  by (rule someI-ex[of λf'. ?is-ext f f']) (smt (verit) that extension-lemma-localE2)
  thus ?is-ext f (extend f) unfolding extend-def by blast
qed

```

```

  have extend ' diffeo-ψ.dest.diff-fun-space ⊆ manifold-eucl.dest.diff-fun-space
  using extend by blast

```

```

  have  $d\kappa^{-1}$  (directional-derivative k (ψ p) x) f = restrict0 (diffeo-ψ.dest.diff-fun-space)
  (λf. frechet-derivative f (at (ψ p)) x) f
  for f
proof (cases f ∈ diffeo-ψ.dest.diff-fun-space)
  case True
  have frechet-derivative-extend: frechet-derivative f (at (ψ p)) x = frechet-derivative
  (extend f) (at (ψ p)) x

```

if $f: f \in \text{diffeo-}\psi.\text{dest.diff-fun-space}$ **for** f
proof –
obtain N **where** $N: \psi p \in N \wedge \text{open } N \wedge \text{closure } N \subseteq \text{diffeo-}\psi.\text{dest.carrier}$
 $\wedge (\forall x \in \text{closure } N. (\text{extend } f) x = f x) \wedge$
 $(\forall v \in T_{\psi p} \psi U. v f = v (\text{extend } f)) \wedge (\forall v \in T_{\psi p} \psi U. v (\text{extend } f) = d\kappa v$
 $(\text{extend } f)) \wedge (\forall v \in T_{\psi p} E. d\kappa^{-1} v f = v (\text{extend } f))$
using $\text{extend}(1)[OF f]$ **by** *presburger*
show *?thesis*
apply (*rule frechet-derivative-transform-within-open-ext*[**where** $f=f$ **and**
 $g=\text{extend } f$ **and** $X=N$ **for** f])
using *sub-eucl.submanifold-atlasI sub-eucl.sub-diff-fun-differentiable-at*
 $[OF \text{diffeo-}\psi.\text{dest.diff-fun-space}D[OF f], \text{of restrict-chart (codomain } \psi$
 $\text{chart-eucl}]$
apply (*simp add: id-def[symmetric] assms*)
using N **by** *simp-all*
qed
have $d\kappa^{-1} (\text{directional-derivative } k (\psi p) x) f = (\text{directional-derivative } k (\psi p)$
 $x) (\text{extend } f)$
using *assms eq-T_{ψp}E-range-inclusion eq-T_{ψp}E-range-inclusion2 extend(1)*
True by blast
also have $\dots = \text{frechet-derivative } (\text{extend } f) (\text{at } (\psi p)) x$
unfolding *directional-derivative-def* **using** $\text{extend}(1)[OF \text{True}]$ **by** *simp*
finally show *?thesis*
using *True frechet-derivative-extend* **by** *simp*
next
case *False*
then show *?thesis*
proof –
have *RHS-0: restrict0 diffeo-ψ.dest.diff-fun-space (λf. frechet-derivative f (at*
 $(\psi p)) x) f = 0$
using *restrict0-apply-out[OF False]* **by** *blast*
moreover have *LHS-0: dκ⁻¹ (directional-derivative k (ψ p) x) f = 0*
using *bij-betwE[OF bij-betw-dκ-inv] bij-betwE[OF bij-betw-directional-derivative[OF*
 $\text{assms}]$
using *diffeo-ψ.dest.tangent-spaceD extensional0-outside[OF False]* **by** *blast*
ultimately show *?thesis* **by** *simp*
qed
qed
thus *?thesis* **by** *blast*
qed

lemma *smooth-on-compat-charts-TM:*

assumes $c1 \in \text{atlas } c2 \in \text{atlas}$
shows *smooth-on* $(c1 \text{ ‘ } (\text{domain } c1 \cap \text{domain } c2) \times \text{UNIV})$
 $(\lambda x. \text{frechet-derivative } ((\lambda y. (\text{restrict-chart } (\text{domain } c1 \cap \text{domain } c2) c2) y \cdot$
 $i) \circ \text{inv-chart } (\text{restrict-chart } (\text{domain } c1 \cap \text{domain } c2) c1)) (\text{at } (\text{fst } x)) (\text{snd } x))$

```

  (is ‹smooth-on ?D (λx. frechet-derivative ((λy. ?r2 y · i) ∘ ?r1i) (at (fst x))
(snd x)))›)
proof –
  let ?dom-Int = domain c1 ∩ domain c2
  have open-simps[simp]: open ?dom-Int open ?D
    by (auto simp: open-Int open-Times)

  have smooth-on-1: smooth-on (fst' ?D) ((λy. ?r2 y · i) ∘ ?r1i) for i
    apply simp
    apply (rule smooth-on-cong'[of - c1 ' (domain c1 ∩ domain c2)])
    apply (rule smooth-on-cong[of - - (λy. c2 (inv-chart c1 y) · i)])
    apply (rule smooth-on-inner[OF - smooth-on-const[of - - i]])
    using atlas-is-atlas[unfolded smooth-compat-def o-def, OF assms(1,2)] apply
  auto[4]
  unfolding restrict-codomain-TM-def codomain-TM-alt using image-domain-eq
  by fastforce
  have smooth-on-2: smooth-on ?D (λx. frechet-derivative ((λy. (?r2 y) · i) ∘ ?r1i)
(at (fst x)) v) for v i
    apply (rule smooth-on-compose2[OF derivative-is-smooth, unfolded o-def, where
S=UNIV and T=fst' ?D])
    using smooth-on-fst smooth-on-1 by (auto simp: open-image-fst)

  have r2-r1i-differentiable: (λx. ?r2 (?r1i x) · i) differentiable (at (fst p)) if p ∈
?D for i::'b and p
  proof –
    have 1: open (c1 ' (domain c1 ∩ domain c2))
      and 2: c1 (inv-chart c1 (fst p)) = fst p
      and 3: inv-chart c1 (fst p) ∈ ?dom-Int using that by auto
    show ?thesis
      using smooth-on-imp-differentiable-on[unfolded differentiable-on-def, OF
smooth-on-1]
      by (simp add: o-def) (metis at-within-open image-eqI 1 2 3)
  qed

  show ?thesis
    unfolding o-def
    apply (rule smooth-on-cong[OF - - frechet-derivative-componentwise[OF r2-r1i-differentiable]])
    apply (rule smooth-on-sum)
    apply (rule smooth-on-times-fun[of ∞ ?D, unfolded times-fun-def])
    subgoal by (auto intro!: smooth-on-inner smooth-on-snd)
    subgoal using smooth-on-2[unfolded o-def] by simp
    by simp-all
  qed

```

— The charts defined above for the tangent bundle of an infinitely smooth manifold are compatible (see *smooth-compat*) if the charts used for the construction are compatible. Thus, we can construct an atlas (up to type class issues) for *TM* from the atlas of the manifold.

lemma *atlas-TM*:

assumes $c1 \in \text{atlas } c2 \in \text{atlas}$

shows $\text{smooth-on } ((\text{apply-chart-TM } c1) \text{ ' } (\text{domain-TM } c1 \cap \text{domain-TM } c2))$
 $((\text{apply-chart-TM } c2) \circ (\text{inv-chart-TM } c1))$
 $(\text{is } \langle \text{smooth-on } (?c1 \text{ ' } (?dom1 \cap ?dom2)) ((?c2) \circ (?i1)) \rangle)$

proof –

let $?dom\text{-}Int = \text{domain } c1 \cap \text{domain } c2$

have $dom\text{-}eq: ?dom1 \cap ?dom2 = \{(p,v). p \in \text{domain } c1 \wedge p \in \text{domain } c2 \wedge v \in \text{tangent-space } p\}$

unfolding *domain-TM-def* **by** *auto*

have $open\text{-}Int\text{-}dom[simp]: \text{open } (\text{domain } c1 \cap \text{domain } c2)$ **by** *blast*

have $open\text{-}image\text{-}dom\text{-}TM[simp]: \text{open } (\text{apply-chart-TM } c1 \text{ ' } (\text{domain-TM } c1 \cap \text{domain-TM } c2))$

using *assms open-intersection-TM* **by** *blast*

have $inv\text{-}chart\text{-}x\text{-}in: (\text{inv-chart } c1 \ x) \in \text{domain } c1 \cap \text{domain } c2$

if $x \in c1 \text{ ' } (\text{domain } c1 \cap \text{domain } c2)$ **for** x

using *that by force*

let $?snd\text{-}c2i1 = \lambda(p, v). c\text{-manifold-point.tangent-chart-fun charts } \infty \ c2 \ (\text{inv-chart } c1 \ p)$

$(c\text{-manifold-point.coordinate-vector charts } \infty \ c1 \ (\text{inv-chart } c1 \ p) \ v)$

let $?R1i = \text{restrict-inv-chart-TM } (\text{domain } c1 \cap \text{domain } c2) \ c1$

and $?R1 = \text{restrict-chart-TM } (\text{domain } c1 \cap \text{domain } c2) \ c1$

and $?R2 = \text{restrict-chart-TM } (\text{domain } c1 \cap \text{domain } c2) \ c2$

and $?r1 = \text{restrict-chart } ?dom\text{-}Int \ c1$

and $?r2 = \text{restrict-chart } ?dom\text{-}Int \ c2$

and $?r1i = \text{inv-chart } (\text{restrict-chart } ?dom\text{-}Int \ c1)$

and $?r2i = \text{inv-chart } (\text{restrict-chart } ?dom\text{-}Int \ c2)$

show *?thesis*

proof (*subst restrict-chart-domain-Int[OF assms(1)], subst image-restrict-domain-TM-eq[OF assms(1)], rule smooth-on-cong*)

fix x **assume** $x: x \in \text{restrict-codomain-TM } (\text{domain } c1 \cap \text{domain } c2) \ c1$

then have $y: (?R1i \ x) \in \text{restrict-domain-TM } (\text{domain } c1 \cap \text{domain } c2) \ c2$

using *inv-image-restrict-codomain-TM-eq[OF assms(1)]*

using *restrict-domain-TM-intersection restrict-domain-TM-intersection'*

by *blast*

show $(\text{apply-chart-TM } c2 \circ \text{inv-chart-TM } c1) \ x = (?R2 \circ ?R1i) \ x$

using *inverse-restrict-chart-TM apply-restrict-chart-TM open-Int-dom x y assms(1,2)* **by** *simp*

next

show $open\text{-}restrict\text{-}codomain[simp]: \text{open } (\text{restrict-codomain-TM } (\text{domain } c1 \cap \text{domain } c2) \ c1)$

by (*simp add: image-restrict-domain-TM-eq[OF assms(1), symmetric] restrict-chart-domain-Int[OF assms(1), symmetric]*)

show $\text{smooth-on } (\text{restrict-codomain-TM } (\text{domain } c1 \cap \text{domain } c2) \ c1) \ (?R2 \circ$

?R1i)
proof (*rule smooth-on-Pair'*[*OF open-restrict-codomain*])
have *fst-eq*: $\text{fst} \circ (?R2 \circ ?R1i) = ?r2 \circ ?r1i \circ \text{fst}$
unfolding *restrict-chart-TM-def restrict-inv-chart-TM-def apply-chart-TM-def inv-chart-TM-def* **by** *auto*
show *smooth-on* (*restrict-codomain-TM* ($\text{domain } c1 \cap \text{domain } c2$) *c1*) ($\text{fst} \circ (?R2 \circ ?R1i)$)
apply (*simp add: fst-eq*)
apply (*rule smooth-on-compose*[*of - c1 ' (domain c1 ∩ domain c2)*])
subgoal using *atlas-is-atlas assms smooth-compat-D1* **by** *blast*
subgoal by (*auto intro: smooth-on-fst*)
subgoal by *simp*
subgoal by *simp*
subgoal unfolding *restrict-codomain-TM-def codomain-TM-alt* **using** *image-domain-eq* **by** *fastforce*
done

let $?g = \lambda x. (\sum_{i \in \text{Basis}.} (\text{frechet-derivative } ((\lambda y. (?r2 y) \cdot i) \circ ?r1i) \text{ (at (fst } x)) \text{ (snd } x)) *_{\mathbb{R}} i$)

have *local-simps*: $?r2 \circ ?r1i = (\lambda x. ?r2 (?r1i x))$
and [*simp*]: $\text{domain } c2 \cap (\text{domain } c1 \cap \text{domain } c2) = \text{domain } c1 \cap \text{domain } c2$
c2
by *auto*
have *r2-r1i-differentiable*: $(\lambda x. ?r2 (?r1i x) \cdot i)$ *differentiable* (*at* ($?r1 p$)) **if** $p \in ?\text{dom-Int}$ **for** $i::'b$ **and** p
apply (*rule differentiable-compose*[*of* $\lambda x. x \cdot i$], *simp*)
apply (*subst local-simps(1)[symmetric]*)
apply (*rule c-manifold.diff-fun-differentiable-at*[*of charts-submanifold ?dom-Int ∞*])
subgoal using *atlas-is-atlas charts-submanifold-def in-charts-in-atlas restrict-chart-in-atlas* **by** *unfold-locales (auto)*
subgoal unfolding *diff-fun-def* **using** *diff-apply-chart*[*of ?r2*] *assms(2)* *restrict-chart-in-atlas* **by** *simp*
subgoal using *restrict-chart-in-atlas*[*OF assms(1)*] *c-manifold-local.sub-ψ*
by (*metis c-manifold-point.axioms(1)*[*OF c-manifold-point*] *domain-restrict-chart inf.left-idem open-Int-dom that*)
using *that* **by** *auto*
have *r2p-deriv*: $\text{frechet-derivative } (\lambda x. - (?r2 p) \cdot i)$ (*at* ($?r1 p$)) = 0 **for** $i::'b$ **and** p **by** *auto*
hence *r2p-differentiable*: $(\lambda x. - (?r2 p) \cdot i)$ *differentiable* (*at* ($?r1 p$)) **for** $i::'b$ **and** p **by** *simp*

show *smooth-on* (*restrict-codomain-TM* ($\text{domain } c1 \cap \text{domain } c2$) *c1*) ($\text{snd} \circ (?R2 \circ ?R1i)$)
proof (*rule smooth-on-cong*[*of - - ?g, OF - open-restrict-codomain*])
fix x **assume** $x: x \in \text{restrict-codomain-TM } (\text{domain } c1 \cap \text{domain } c2) \text{ } c1$
then obtain $x_p \ x_v$ **where** *Pair-x*: $x = (x_p, x_v)$ **and** $x_p: x_p \in \text{codomain } c1$ $x_v \in \text{inv-chart } c1 -' ?\text{dom-Int}$

```

unfolding restrict-codomain-TM-def codomain-TM-alt
using codomain-restrict-chart[OF open-Int-dom, of c1] by blast

obtain p where p-def: p = inv-chart ?r1 x_p and p[simp]: p ∈ ?dom-Int
using x_p(2) by auto

interpret p1: c-manifold-point charts ∞ ?r1 p
using x_p(2) by (auto intro!: c-manifold-point simp add: restrict-chart-in-atlas
assms(1) p-def)
interpret p2: c-manifold-point charts ∞ ?r2 p
using x_p(2) by (auto intro!: c-manifold-point simp add: restrict-chart-in-atlas
assms(2) p-def)

let ?v = p1.coordinate-vector x_v
obtain v where v-def: v = p1.coordinate-vector x_v and v[simp]: v ∈
tangent-space p
using p1.coordinate-vector-surj by blast

have pvx: ?R1 (p,v) = x
using Pair-x x_p(1) p1.tangent-chart-fun-inverse(2)
by (auto simp: p-def v-def restrict-chart-TM-def apply-chart-TM-def)

have p1-coord-in-Tp2M: p1.coordinate-vector x_v ∈ p2.T_p M
using v v-def by auto

have diff-fun-spaces-eq[simp]: p2.sub-ψ.sub.diff-fun-space = p1.sub-ψ.sub.diff-fun-space
unfolding p2.sub-ψ.sub.diff-fun-space-def p1.sub-ψ.sub.diff-fun-space-def
by simp
have TpU-eq[simp]: p2.T_p U = p1.T_p U
unfolding p2.sub-ψ.sub.tangent-space-def p1.sub-ψ.sub.tangent-space-def
by simp
have sub-carriers-eq[simp]: p2.sub-ψ.sub.carrier = p1.sub-ψ.sub.carrier
unfolding p2.sub-ψ.sub.carrier-def p1.sub-ψ.sub.carrier-def by simp

have in-diff-fun-space: restrict0 ?dom-Int (λx. (?r2 x - ?r2 p) • i) ∈
p1.sub-ψ.sub.diff-fun-space
for i::'b
proof —
have diff-fun ∞ (charts-submanifold ?dom-Int) (λx. (?r2 x - ?r2 p) • i)
proof (rule diff-fun.diff-fun-cong)
show diff-fun ∞ (charts-submanifold ?dom-Int) ((λx. x • i) ∘ ((λx. (x -
?r2 p)) ∘ ?r2))
proof (intro diff-fun-compose diff-compose)
— This result could easily be an instance of an axiom of Lie groups. However,
I think it may be harder to start from differentiability of a binary operation on the
product manifold than it is to just use composition of basic smooth operations.
have eucl-diff-add-uminus: diff ∞ charts-eucl charts-eucl (λy. y + - x)
if x: x ∈ manifold-eucl.carrier for x::'b
apply (intro diff-fun-charts-euclI[unfolded diff-fun-def])

```

```

    using smooth-on-add[OF smooth-on-id smooth-on-const[of  $\infty$  UNIV
-x]] open-UNIV by simp
    show diff  $\infty$  (charts-submanifold ?dom-Int) (manifold-eucl.dest.charts-submanifold
(codomain ?r2)) ?r2
    using p2.diffeo- $\psi$ .diff-axioms by auto
    show diff  $\infty$  (manifold-eucl.dest.charts-submanifold (codomain ?r2))
charts-eucl ( $\lambda x. x - ?r2 p$ )
    using eucl-diff-add-uminus[of ?r2 p] diff.diff-submanifold p2.sub-eucl.open-submanifold
by auto
    show diff-fun  $\infty$  charts-eucl ( $\lambda x. x \cdot i$ )
    using smooth-on-inner-const by (simp add: diff-fun-charts-eucl)
    qed
    qed (simp)
    moreover have (?r2 x - ?r2 p)  $\cdot$  i = (restrict0 ?dom-Int ( $\lambda x. (?r2 x -$ 
?r2 p)  $\cdot$  i)) x
    if  $x \in$  manifold.carrier (charts-submanifold ?dom-Int) for x
    using p1.sub- $\psi$ -carrier that by auto
    ultimately show ?thesis
    using p1.sub- $\psi$ .sub.restrict0-in-fun-space p2.sub- $\psi$ -carrier by auto
    qed

have p2-comp-p1-coord- $x_v$ : p2.component-function (p1.coordinate-vector  $x_v$ )
i =
    frechet-derivative (( $\lambda y. (?r2 y) \cdot i$ )  $\circ$  ?r1i) (at (?r1 p))  $x_v$  for i::'b
    proof -
    have 1: p2.component-function (p1.coordinate-vector  $x_v$ ) i =
        (p1.differential-inv-chart (p1.dRestr (directional-derivative  $\infty$  (?r1 p)
 $x_v$ ))) (restrict0 ?dom-Int ( $\lambda x. (?r2 x - ?r2 p) \cdot i$ ))
    proof -
    have p2.component-function (p1.coordinate-vector  $x_v$ ) i =
        p2.dRestr2 (p1.coordinate-vector  $x_v$ ) (restrict0 ?dom-Int ( $\lambda x. (?r2 x$ 
- ?r2 p)  $\cdot$  i))
    using p2.component-function-apply-in- $T_p M$ [OF p1-coord-in- $Tp2M$ ]
    by (simp add: Int-absorb1)
    also have ... = p1.dRestr2 (p1.coordinate-vector  $x_v$ ) (restrict0 ?dom-Int
( $\lambda x. (?r2 x - ?r2 p) \cdot i$ ))
    unfolding the-inv-into-def by simp
    also have ... = (p1.differential-inv-chart (p1.dRestr (directional-derivative
 $\infty$  (?r1 p)  $x_v$ )))
        (restrict0 ?dom-Int ( $\lambda x. (?r2 x - ?r2 p) \cdot i$ ))
    using the-inv-into-f-f[OF bij-betw-imp-inj-on[OF p1.tangent-submanifold-isomorphism(1)]]
    using bij-betwE[OF p1.bij-betw-d $\psi$ -inv] bij-betwE[OF p1.bij-betw-d $\kappa$ -inv]
p1-coord-in- $Tp2M$ 
    by (auto simp: p1.coordinate-vector-apply)
    finally show ?thesis .
    qed
    also have ... = frechet-derivative (restrict0 p1.diffeo- $\psi$ .dest.carrier
((restrict0 ?dom-Int ( $\lambda x. (?r2 x - ?r2 p) \cdot i$ )  $\circ$  ?r1i)) (at (?r1 p))  $x_v$ )
    proof -

```

have ... = (p1.differential-inv-chart (restrict0 (p1.diffeo-ψ.dest.diff-fun-space)
(λf. frechet-derivative f (at (?r1 p)) x_v)))
(restrict0 ?dom-Int (λx. (?r2 x - ?r2 p) · i))
using p1.dκ-inv-directional-derivative-eq **by** simp
also have ... = (λg. restrict0 p1.diffeo-ψ.dest.diff-fun-space (λf.
frechet-derivative f (at (?r1 p)) x_v)
(restrict0 p1.diffeo-ψ.dest.carrier (g ∘ ?r1i)))
(restrict0 ?dom-Int (λx. (?r2 x - ?r2 p) · i))
unfolding p1.diffeo-ψ.inv.push-forward-def **using** in-diff-fun-space **by**
simp
also have ... = (λf. frechet-derivative f (at (?r1 p)) x_v)
(restrict0 p1.diffeo-ψ.dest.carrier ((restrict0 ?dom-Int (λx. (?r2 x -
?r2 p) · i)) ∘ ?r1i))
using in-diff-fun-space p1.diffeo-ψ.inv.restrict-compose-in-diff-fun-space
by auto
finally show ?thesis **by** simp
qed
also have ... = frechet-derivative ((λx. (?r2 x) · i) ∘ ?r1i) (at (?r1 p)) x_v
proof -
let ?X = p1.diffeo-ψ.dest.carrier
have X-eq-codomain-r1[simp]: p1.diffeo-ψ.dest.carrier = codomain ?r1
using chart-eucl-simps(1) manifold.carrier-def
by (metis (no-types, lifting) Int-UNIV-right Int-commute ccpo-Sup-singleton
image-insert image-is-empty p1.sub-eucl.carrier-submanifold)
have 1: frechet-derivative (restrict0 p1.diffeo-ψ.dest.carrier ((restrict0
?dom-Int (λx. (?r2 x - ?r2 p) · i) ∘ ?r1i)) (at (?r1 p))) =
frechet-derivative ((λx. (?r2 x - ?r2 p) · i) ∘ ?r1i) (at (?r1 p))
(is ⟨frechet-derivative ?f_L (at -) = frechet-derivative ?f_R (at -)⟩)
proof (rule frechet-derivative-transform-within-open)
show ?f_L x = ?f_R x **if** x ∈ ?X **for** x
using X-eq-codomain-r1 **that** **by** simp
show open ?X **by** blast
show ?r1 p ∈ ?X **using** p1.ψp-in **by** blast
let ?f_L' = (restrict0 ?dom-Int (λx. (?r2 x - ?r2 p) · i) ∘ ?r1i
show ?f_L differentiable at (?r1 p)
apply (rule differentiable-transform-within-open[of ?f_L' - - ?X])
apply (rule p1.sub-ψ.sub-diff-fun-differentiable-at)
using p1.ψp-in p1.diffeo-ψ.dest.open-carrier in-diff-fun-space p1.sub-ψ
p1.p p1.sub-ψ.sub.diff-fun-spaceD **by** auto
qed
also have 2: ... = frechet-derivative ((λx. (?r2 x) · i) ∘ ?r1i) (at (?r1
p))
proof -
have frechet-derivative ((λx. (?r2 x - ?r2 p) · i) ∘ ?r1i) (at (?r1 p)) =
frechet-derivative ((λx. ?r2 (?r1i x) · i) + (λx. - (?r2 p) · i))
(at (?r1 p))
by (simp add: plus-fun-def inner-diff-left) (meson comp-apply)
also have ... = frechet-derivative (λx. ?r2 (?r1i x) · i) (at (?r1 p)) +
(frechet-derivative (λx. - (?r2 p) · i) (at (?r1 p)))


```

      using r2-r1i-differentiable[OF p] r2p-differentiable by (auto simp:
frechet-derivative-plus-fun)
      finally show frechet-derivative ((λx. (?r2 x - ?r2 p) · i) ∘ ?r1i) (at
(?r1 p)) =
        frechet-derivative ((λx. ?r2 x · i) ∘ ?r1i) (at (?r1 p))
      using r2p-deriv by simp (metis comp-apply)
    qed
    finally show ?thesis using 1 by presburger
  qed
  finally show p2.component-function (p1.coordinate-vector xv) i =
frechet-derivative ((λx. (?r2 x) · i) ∘ ?r1i) (at (?r1 p)) xv .
  qed

  have (snd ∘ (?R2 ∘ ?R1i)) x = (p2.tangent-chart-fun ∘ p1.coordinate-vector)
xv
    unfolding restrict-chart-TM-def restrict-inv-chart-TM-def
    unfolding apply-chart-TM-def inv-chart-TM-def by (simp add: Pair-x
p-def)
  then show (snd ∘ (?R2 ∘ ?R1i)) x = ?g x
    unfolding p2.tangent-chart-fun-def using p2-comp-p1-coord-xv p-def p
Pair-x xp(1) by auto
  next
  let ?D = restrict-codomain-TM (domain c1 ∩ domain c2) c1
  show smooth-on ?D ?g
  proof (rule smooth-on-sum, rule smooth-on-scaleR)
    fix i::'b assume i: i∈Basis
    have D-product: ?D = c1 ‘ (domain c1 ∩ domain c2) × UNIV
      unfolding restrict-codomain-TM-def codomain-TM-def
      by auto (metis IntI chart-inverse-inv-chart imageI inv-chart-in-domain)
    show smooth-on ?D (λx. frechet-derivative ((λy. (?r2 y) · i) ∘ ?r1i) (at
(fst x)) (snd x))
      unfolding D-product by (rule smooth-on-compat-charts-TM[OF assms])
    qed (auto)
  qed
  qed
  qed
  qed
  qed

```

lemma atlas-TM':
assumes $c1 \in \text{atlas}$ $c2 \in \text{atlas}$
shows $\text{smooth-on } ((\text{apply-chart-TM } c2) ‘ (\text{domain-TM } c1 \cap \text{domain-TM } c2))$
 $((\text{apply-chart-TM } c1) \circ (\text{inv-chart-TM } c2))$
using $\text{atlas-TM}[OF \text{assms}(2,1)]$ **by** (simp add: Int-commute)

end

5.1.3 Differentiability of vector fields

context $c\text{-manifold}$ **begin**

abbreviation $k\text{-diff-from-}M\text{-to-}TM\text{-at-in} :: \text{enat} \Rightarrow 'a \Rightarrow ('a, 'b)\text{chart} \Rightarrow ('a \Rightarrow 'a \text{ tangent-bundle}) \Rightarrow \text{bool}$

where $k\text{-diff-from-}M\text{-to-}TM\text{-at-in } k' x c X \equiv x \in \text{domain } c \wedge X \text{ ' domain } c \subseteq \text{domain-}TM \ c \wedge k'\text{-smooth-on } (\text{codomain } c) \ (\text{apply-chart-}TM \ c \circ X \circ \text{inv-chart } c)$

— Compare this definition to $\text{diff-axioms } ?k \ ?charts1.0 \ ?charts2.0 \ ?f \equiv \forall x. x \in \text{manifold.carrier } ?charts1.0 \longrightarrow (\exists c1 \in c\text{-manifold.atlas } ?charts1.0 \ ?k. \exists c2 \in c\text{-manifold.atlas } ?charts2.0 \ ?k. x \in \text{domain } c1 \wedge ?f \text{ ' domain } c1 \subseteq \text{domain } c2 \wedge ?k\text{-smooth-on } (\text{codomain } c1) \ (\text{apply-chart } c2 \circ ?f \circ \text{inv-chart } c1))$. It's the same, except the charts for TM aren't of type $('a, 'b)\text{chart}$.

definition $k\text{-diff-from-}M\text{-to-}TM \ (\langle \text{--diff'--from-}M'\text{-to-}TM \rangle [1000])$

where $\text{diff-from-}M\text{-to-}TM\text{-def}: k'\text{-diff-from-}M\text{-to-}TM \ X \equiv \forall x. x \in \text{carrier} \longrightarrow (\exists c \in \text{atlas}. k'\text{-diff-from-}M\text{-to-}TM\text{-at-in } k' x c X)$

abbreviation $\text{continuous-from-}M\text{-to-}TM \equiv 0\text{-diff-from-}M\text{-to-}TM$

abbreviation (in smooth-manifold) $\text{smooth-from-}M\text{-to-}TM \equiv k\text{-diff-from-}M\text{-to-}TM \ \infty$

lemma $\text{diff-from-}M\text{-to-}TM\text{-E}$:

assumes $k'\text{-diff-from-}M\text{-to-}TM \ X \ x \in \text{carrier}$

obtains c **where** $c \in \text{atlas} \ x \in \text{domain } c \ X \text{ ' domain } c \subseteq \text{domain-}TM \ c \ k'\text{-smooth-on } (\text{codomain } c) \ (\text{apply-chart-}TM \ c \circ X \circ \text{inv-chart } c)$

using *assms* **unfolding** $\text{diff-from-}M\text{-to-}TM\text{-def}$ **by** *auto*

lemma $\text{continuous-from-}M\text{-to-}TM\text{-D}$:

assumes $\text{continuous-from-}M\text{-to-}TM \ X \ x \in \text{carrier}$

obtains c **where** $c \in \text{atlas} \ x \in \text{domain } c \ X \text{ ' domain } c \subseteq \text{domain-}TM \ c \ \text{continuous-on } (\text{codomain } c) \ (\text{apply-chart-}TM \ c \circ X \circ \text{inv-chart } c)$

using *assms* **by** (meson $\text{diff-from-}M\text{-to-}TM\text{-E}$ $\text{smooth-on-imp-continuous-on}$ that)

definition $\text{section-of-}TM\text{-def}: \text{section-of-}TM\text{-on } S \ X \equiv \forall p \in S. (X \ p) \in TM \wedge \text{fst } (X \ p) = p$

abbreviation $\text{section-of-}TM \equiv \text{section-of-}TM\text{-on } \text{carrier}$

lemma $\text{section-of-}TM\text{-subset}$:

assumes $\text{section-of-}TM\text{-on } S \ X \ T \subseteq S$

shows $\text{section-of-}TM\text{-on } T \ X$

using *assms* **unfolding** $\text{section-of-}TM\text{-def}$ **by** *force*

lemma $\text{section-domain-}TM$:

assumes $\text{section-of-}TM\text{-on } (\text{domain } c) \ X$

shows $X \text{ ' domain } c \subseteq \text{domain-}TM \ c$

using *assms* **unfolding** $\text{domain-}TM\text{-def}$ $\text{section-of-}TM\text{-def}$ $\text{in-}TM\text{-def}$ **by** *auto*

lemma $\text{section-domain-}TM'$:

assumes $\text{section-of-}TM \ X \ c \in \text{atlas}$

shows $X \text{ ' domain } c \subseteq \text{domain-}TM \ c$

using *assms section-domain-TM section-of-TM-subset* **by** *blast*

lemma *section-vimage-domain-TM*:

assumes *section-of-TM X c ∈ atlas*

shows $\text{carrier} \cap X - ' \text{domain-TM } c = \text{domain } c$

using *assms unfolding domain-TM-def section-of-TM-def in-TM-def*

by *simp force*

end

context *smooth-manifold begin*

Show that a smooth/differentiable vector field is smooth in any chart. This would be $\llbracket \text{diff } ?k \text{ ?charts1.0 ?charts2.0 ?f; ?d1.0} \in \text{c-manifold.atlas ?charts1.0 ?k; ?d2.0} \in \text{c-manifold.atlas ?charts2.0 ?k} \rrbracket \implies ?k\text{-smooth-on } (\text{codomain } ?d1.0 \cap \text{inv-chart } ?d1.0 - ' (\text{manifold.carrier } ?charts1.0 \cap ?f - ' \text{domain } ?d2.0)) (\text{apply-chart } ?d2.0 \circ ?f \circ \text{inv-chart } ?d1.0)$ if we could write TM as a c -manifold; it relies on the compatibility of charts for TM given in $\llbracket \text{smooth-manifold } ?charts; ?c1.0} \in \text{c-manifold.atlas ?charts} \infty; ?c2.0} \in \text{c-manifold.atlas ?charts} \infty \rrbracket \implies \text{smooth-on } (\text{c-manifold.apply-chart-TM } ?charts ?c1.0 - ' (\text{c-manifold.domain-TM } ?charts \infty ?c1.0 \cap \text{c-manifold.domain-TM } ?charts \infty ?c2.0)) (\text{c-manifold.apply-chart-TM } ?charts ?c2.0 \circ \text{c-manifold.inv-chart-TM } ?charts ?c1.0)$.

lemma *diff-from-M-to-TM-chartsD*:

assumes $X: k\text{-diff-from-M-to-TM } k' X$ **section-of-TM** X **and** $c: c \in \text{atlas}$

shows $k'\text{-smooth-on } (\text{codomain } c) (\text{apply-chart-TM } c \circ X \circ \text{inv-chart } c)$

proof –

have *codom-simp*: $\text{codomain } c \cap \text{inv-chart } c - ' (\text{carrier} \cap X - ' \text{domain-TM } c) = \text{codomain } c$

using *section-vimage-domain-TM[OF X(2) c]* **by** (*simp add: Int-absorb2 subset-vimage-iff*)

{ **fix** y **assume** $y \in \text{codomain } c \cap \text{inv-chart } c - ' (\text{carrier} \cap X - ' \text{domain-TM } c)$

then have $y: X (\text{inv-chart } c y) \in \text{domain-TM } c y \in \text{codomain } c$

by *auto*

then obtain x **where** $x: c x = y x \in \text{domain } c$

by *force*

then have $x \in \text{carrier}$ **using** *assms* **by** *force*

obtain $c1$ **where** $c1 \in \text{atlas}$

and $fc1: X - ' \text{domain } c1 \subseteq \text{domain-TM } c1$

and $xc1: x \in \text{domain } c1$

and $d: k'\text{-smooth-on } (\text{codomain } c1) (\text{apply-chart-TM } c1 \circ X \circ \text{inv-chart } c1)$

by (*meson* $\langle x \in \text{carrier} \rangle$ *assms(1) diff-from-M-to-TM-E*)

have $fc1' [simp]: x \in \text{domain } c1 \implies X x \in \text{domain-TM } c1$ **for** x **using** $fc1$ **by** *auto*

have $r1: k'\text{-smooth-on } (c - ' (\text{domain } c \cap \text{domain } c1)) (c1 \circ \text{inv-chart } c)$

using *smooth-compat-D1[OF smooth-compat-le[OF atlas-is-atlas[OF c* $\langle c1 \in \text{atlas} \rangle$]]] **by** *force*

— Important: this is where we use $\llbracket \text{smooth-manifold } ?charts; ?c1.0 \in c\text{-manifold.atlas } ?charts \infty; ?c2.0 \in c\text{-manifold.atlas } ?charts \infty \rrbracket \implies \text{smooth-on } (c\text{-manifold.apply-chart-TM } ?charts ?c1.0 \text{ ' } (c\text{-manifold.domain-TM } ?charts \infty ?c1.0 \cap c\text{-manifold.domain-TM } ?charts \infty ?c2.0)) (c\text{-manifold.apply-chart-TM } ?charts ?c2.0 \circ c\text{-manifold.inv-chart-TM } ?charts ?c1.0)$.

have $r2: k'\text{-smooth-on } (\text{apply-chart-TM } c1 \text{ ' } (\text{domain-TM } c \cap \text{domain-TM } c1))$
 $(\text{apply-chart-TM } c \circ \text{inv-chart-TM } c1)$
apply $(\text{rule smooth-on-le}[OF \text{ atlas-TM}'[OF c \langle c1 \in \text{atlas} \rangle]])$ **by** *simp*

define T **where** $T = c \text{ ' } (\text{domain } c \cap \text{domain } c1) \cap \text{inv-chart } c \text{ - ' } (\text{carrier} \cap (X \text{ - ' } \text{domain-TM } c))$

have $\text{simps-1}: (\text{apply-chart-TM } c1 \circ X \circ \text{inv-chart } c1) \text{ ' } (\text{apply-chart } c1 \circ \text{inv-chart } c) \text{ ' } T = (\text{apply-chart-TM } c1 \circ X \circ \text{inv-chart } c) \text{ ' } T$
if $\text{inv-chart } c \text{ ' } T \subseteq \text{domain } c1$ **for** T
unfolding $\text{image-comp}[\text{symmetric}]$ **using** *that by auto*
 $(\text{smt } (\text{verit}) \text{ image-eqI image-subset-iff inv-chart-inverse})$
have $\text{inv-chart } c \text{ ' } T \subseteq \text{domain } c1$
by $(\text{auto simp: } T\text{-def})$
note $T\text{-simps} = \text{simps-1}[OF \text{ this}] \text{ section-vimage-domain-TM}[OF X(2) c]$
have *open* T
by $(\text{auto intro!}: \text{open-continuous-vimage}' \text{ continuous-intros simp: } T\text{-simps}(2) T\text{-def})$

have $T\text{-subset}: T \subseteq \text{apply-chart } c \text{ ' } (\text{domain } c \cap \text{domain } c1)$
by $(\text{auto simp: } T\text{-def})$
have $\text{opens}: \text{open } (c1 \text{ ' } \text{inv-chart } c \text{ ' } T) \text{ open } (\text{apply-chart-TM } c1 \text{ ' } (\text{domain-TM } c \cap \text{domain-TM } c1))$
using $T\text{-subset } fc1 \langle \text{open } T \rangle \langle \text{inv-chart } c \text{ ' } T \subseteq \text{domain } c1 \rangle$ **apply** *blast*
by $(\text{metis Int-commute } \langle c1 \in \text{atlas} \rangle \text{ open-intersection-TM})$
have $k'\text{-smooth-on } ((\text{apply-chart } c1 \circ \text{inv-chart } c) \text{ ' } T) (\text{apply-chart-TM } c \circ \text{inv-chart-TM } c1 \circ (\text{apply-chart-TM } c1 \circ X \circ \text{inv-chart } c1))$
using $r2 d \text{ opens unfolding image-comp}[\text{symmetric}]$ **apply** $(\text{rule smooth-on-compose2})$
by $(\text{auto simp: } T\text{-def}) (\text{metis IntI } fc1 \text{ image-subset-iff subset-refl})$
from $\text{this } r1 \langle \text{open } T \rangle \text{ opens}(1)$ **have** $k'\text{-smooth-on } T$
 $((\text{apply-chart-TM } c \circ \text{inv-chart-TM } c1) \circ (\text{apply-chart-TM } c1 \circ X \circ \text{inv-chart } c1) \circ (c1 \circ \text{inv-chart } c))$
unfolding $\text{image-comp}[\text{symmetric}]$
by $(\text{rule smooth-on-compose2}) (\text{force simp: } T\text{-def})+$
then **have** $k'\text{-smooth-on } T (\text{apply-chart-TM } c \circ X \circ \text{inv-chart } c)$
using $\langle \text{open } T \rangle$ **apply** $(\text{rule smooth-on-cong})$
using $\text{apply-chart-TM-inverse}(1)[\text{of } c1 \text{ fst } (X \text{ } xa) \text{ snd } (X \text{ } xa) \text{ for } xa] fc1' \langle c1 \in \text{atlas} \rangle$
by $(\text{auto simp: } T\text{-def})$
moreover **have** $y \in T$
using $x \text{ } xc1 \text{ } fc1 \text{ } y \langle c1 \in \text{atlas} \rangle$ **by** $(\text{auto simp: } T\text{-def})$
ultimately **have** $\exists T. y \in T \wedge \text{open } T \wedge k'\text{-smooth-on } T (\text{apply-chart-TM } c \circ X \circ \text{inv-chart } c)$
using $\langle \text{open } T \rangle$ **by** *metis* }

thus *?thesis*
apply (*rule smooth-on-open-subsetsI*)
using *codom-simp* **by** *simp*
qed

definition *smooth-section-of-TM* $X \equiv \text{section-of-TM } X \wedge \text{smooth-from-M-to-TM } X$

abbreviation *set-of-smooth-sections-of-TM* $(\langle \mathfrak{X} \rangle)$
where *set-of-smooth-sections-of-TM* $\equiv \{X. \text{smooth-section-of-TM } X\}$

lemma *in \mathfrak{X} -E*:
assumes $X \in \mathfrak{X} \ p \in \text{carrier}$
shows $(\exists c \in \text{atlas}. p \in \text{domain } c \wedge X \text{ 'domain } c \subseteq \text{domain-TM } c \wedge \text{smooth-on}(\text{codomain } c) (\text{apply-chart-TM } c \circ X \circ \text{inv-chart } c))$
and $\text{snd } (X \ p) \in \text{tangent-space } p$
and $\text{fst } (X \ p) = p$
using *assms TM-E[$\text{of } X \ p \ \text{for } p$]*
by (*auto simp: smooth-section-of-TM-def section-of-TM-def diff-from-M-to-TM-def*)
(metis)

lemma *in \mathfrak{X} -chartsD*:
assumes $X \in \mathfrak{X} \ c \in \text{atlas}$
shows *smooth-on* (*codomain* c) (*apply-chart-TM* $c \circ X \circ \text{inv-chart } c$)
using *diff-from-M-to-TM-chartsD[$\text{of } \infty \ X \ c$]* *assms smooth-section-of-TM-def* **by**
auto

end

A vector field is smooth if it is smooth as a map $M \rightarrow TM$. As a shortcut, we define a smooth vector field as one that is smooth in the chart - this avoids problems with defining a $(\text{'a} \times ((\text{'a} \Rightarrow \text{real}) \Rightarrow \text{real}), \text{'b})$ chart. We also introduce a duality of predicates with strongly related meaning: this allows us to consider vector fields as either maps $\text{'a} \Rightarrow (\text{'a} \Rightarrow \text{real}) \Rightarrow \text{real}$, i.e. mapping a point to a vector; or maps $\text{'a} \Rightarrow \text{'a} \times ((\text{'a} \Rightarrow \text{real}) \Rightarrow \text{real})$, i.e. sections of TM properly speaking.

context *c-manifold* **begin**

definition *rough-vector-field* $:: \text{'a vector-field} \Rightarrow \text{bool}$
where *rough-vector-field* $X \equiv \text{extensional0 carrier } X \wedge (\forall p \in \text{carrier}. X \ p \in \text{tangent-space } p)$

lemma *rough-vector-fieldE* [*elim*]:
assumes *rough-vector-field* X
shows $\bigwedge p. X \ p \in \text{tangent-space } p \ \text{extensional0 carrier } X$
using *assms* **by** (*auto simp: rough-vector-field-def extensional0-outside tangent-space.mem-zero*)

lemma *rough-vector-field-subset*:

assumes *rough-vector-field* $X \ T \subseteq \text{carrier}$
shows *rough-vector-field* (*restrict0* $T \ X$)
unfolding *rough-vector-field-def* **using** *assms* *rough-vector-fieldE* *tangent-space.mem-zero*
by (*metis* (*no-types*, *lifting*) *extensional0-def* *restrict0-def*)

end

abbreviation (*input*) *vec-field-apply-fun* :: ' a *vector-field* \Rightarrow (' $a \Rightarrow \text{real}$) \Rightarrow (' $a \Rightarrow \text{real}$)
(**infix** ' \rangle ' 100)
where *vec-field-apply-fun* $X \ f \equiv \lambda p. \ X \ p \ f$

lemma (**in** *c-manifold*) *vec-field-apply-fun-cong*:
assumes X : *rough-vector-field* X **and** U : *open* $U \ U \subseteq \text{carrier} \ \forall x \in U. \ f \ x = g \ x$
and f : $f \in \text{diff-fun-space}$ **and** g : $g \in \text{diff-fun-space}$
shows $\forall p \in U. \ X \ p \ f = X \ p \ g$
using *assms* **by** (*auto* *intro*: *derivation-eq-localI* *simp*: *rough-vector-field-def*)

lemma (**in** *c-manifold*) *ext0-vec-field-apply-fun*:
assumes X : *rough-vector-field* X
shows *extensional0* *diff-fun-space* (*vec-field-apply-fun* X)
using *rough-vector-fieldE*[*OF* X] **unfolding** *tangent-space-def* *extensional0-def*
by *fastforce*

5.2 Smoothness criterion for a vector field in a single chart.

A smooth vector field is one that is infinitely differentiable when expanded in the charting Euclidean space using $\llbracket c\text{-manifold-point } ?charts \ ?k \ ?\psi \ ?p; \ ?v \in c\text{-manifold.tangent-space } ?charts \ ?k \ ?p; \ ?k = \infty \rrbracket \Longrightarrow \ ?v = (\sum_{i \in \text{Basis. } c\text{-manifold-point.component-function } ?charts \ ?k \ ?\psi \ ?p} \ ?v \ i \ *_{\mathbb{R}} \ c\text{-manifold-point.coordinate-vector } ?charts \ ?k \ ?\psi \ ?p \ i)$. This should be the chart that makes each tangent space into a manifold anyway, but the type constraints are tricky to satisfy.

Since tangent spaces at the same point differ between a manifold and a submanifold, it's important to note that the differentiability condition can be relaxed to only apply to a subset, but the tangent bundle is always the disjoint union of tangent spaces of the *entire* manifold, which implies the chart function for the tangent space is defined in the entire manifold, not a submanifold.

locale *smooth-vector-field-local* = *c-manifold-local* *charts* ∞ ψ **for** *charts* ψ +
fixes X
assumes *vector-field*: $\forall p \in \text{domain } \psi. \ X \ p \in \text{tangent-space } p$
and *smooth-in-chart*: *diff-fun* ∞ (*charts-submanifold* (*domain* ψ)) ($\lambda p. (c\text{-manifold-point.tangent-chart-fun } \text{charts } \infty \ \psi \ p) \ (X \ p)$)
begin
lemma *rough-vector-field*: *rough-vector-field* (*restrict0* (*domain* ψ) X)
apply (*simp* *only*: *rough-vector-field-def*, *intro* *conjI*)

```

using extensional0-def sub-ψ-carrier apply fastforce
using vector-field by (metis restrict0-apply-in restrict0-apply-out tangent-space.mem-zero)
end

```

5.2.1 Connecting the types $'a \Rightarrow ('a \Rightarrow \text{real}) \Rightarrow \text{real}$ (used for *smooth-vector-field-local*) and $'a \Rightarrow 'a \times (('a \Rightarrow \text{real}) \Rightarrow \text{real})$ (used for $\lambda \text{charts } k. c\text{-manifold.section-of-TM-on charts } k$ (*manifold.carrier charts*)).

context *c-manifold* **begin**

```

lemma fst-apply-chart-TM-id [simp]: (fst ∘ (apply-chart-TM ψ ∘ X ∘ inv-chart
ψ)) x = x
if section-of-TM-on (domain ψ) X ψ ∈ atlas x ∈ codomain ψ for x
using that by (simp add: case-prod-beta' apply-chart-TM-def section-of-TM-def)

```

The justification for the definition of *smooth-vector-field-local* is the lemma below, connecting it to the smoothness requirement used to define the set of smooth sections \mathfrak{X} .

```

lemma apply-chart-TM-chartX:
fixes X :: ('a ⇒ 'a × (('a ⇒ real) ⇒ real)) and c :: ('a, 'b) chart and chart-X
:: 'a ⇒ 'b
defines chart-X ≡ λp. (c-manifold-point.tangent-chart-fun charts ∞ c p) (snd
(X p))
assumes k: k=∞ and X: section-of-TM-on (domain c) X and c: c ∈ atlas
shows smooth-on (codomain c) (apply-chart-TM c ∘ X ∘ inv-chart c) ⟷ diff-fun
∞ (charts-submanifold (domain c)) chart-X
(is ⟨?smooth-in-chart-TM c X ⟷ ?diff-domain c chart-X⟩)

```

proof –

```

interpret c: c-manifold-local charts ∞ c
using k c pairwise-compat by unfold-locales simp-all
have p: c-manifold-point charts ∞ c p if p ∈ domain c for p
using that by unfold-locales simp
have X-in-TM: fst (X p) = p snd (X p) ∈ tangent-space p if p ∈ domain c for p
using that X c.ψ in-TM-E(1) by (auto simp: section-of-TM-def)
have chart-X-alt: chart-X p = (snd ∘ (c.apply-chart-TM c ∘ X)) p if p ∈ domain
c for p
by (simp add: that chart-X-def c.apply-chart-TM-def X-in-TM(1) split-beta)
have smooth-comp-snd: smooth-on (codomain c) (snd ∘ f) if smooth-on (codomain
c) f for f :: 'b ⇒ 'b × 'b
using open-codomain that by (auto intro!: smooth-on-snd simp: comp-def)
have c-in-sub-atlas: c ∈ c.sub-ψ.sub.atlas
by (metis c.ψ c.atlas-is-atlas c.sub-ψ.sub.maximal-atlas c.sub-ψ.submanifold-atlasE
c.sub-ψ-carrier set-eq-subset)

```

show ?thesis

proof

```

assume asm: ?smooth-in-chart-TM c X

```

```

    have 1: smooth-on (codomain c) (snd ∘ (c.apply-chart-TM c ∘ X ∘ inv-chart
c))
      using smooth-comp-snd[OF asm] by (simp only: comp-assoc)
    have 2: smooth-on (codomain c) ((λx. x) ∘ chart-X ∘ inv-chart c)
      by (auto intro!: smooth-on-cong[OF 1] simp: chart-X-alt)
  {
    fix x assume x ∈ domain c
    then interpret x: c-manifold-point charts ∞ c x using p by blast
    have ∃ c1 ∈ c.sub-ψ.sub.atlas. ∃ c2 ∈ manifold-eucl.atlas ∞.
      x ∈ domain c1 ∧ chart-X ' domain c1 ⊆ domain c2 ∧
      smooth-on (codomain c1) (apply-chart c2 ∘ chart-X ∘ inv-chart c1)
      using 2 c-in-sub-atlas by (intro bexI) auto
  }

  then show ?diff-domain c chart-X
    unfolding diff-fun-def diff-def diff-axioms-def
    using c.sub-ψ.sub.manifold-eucl.c-manifolds-axioms c.sub-ψ-carrier by blast
  next

  assume asm: ?diff-domain c chart-X
  interpret asm-df: diff-fun ∞ charts-submanifold (domain c) snd ∘ (c.apply-chart-TM
c ∘ X)
    using diff-fun.diff-fun-cong[OF asm chart-X-alt] by fastforce
  have codomain-c-eq: codomain c = codomain c ∩ inv-chart c - ' (c.sub-ψ.sub.carrier
∩ (snd ∘ (c.apply-chart-TM c ∘ X)) - ' domain chart-eucl)
    using c.ψ by (simp, blast)
  let ?X = (c.apply-chart-TM c ∘ X ∘ inv-chart c)
  let ?X' = λx. (x, (snd ∘ ?X) x)
  have X-eq: ?X x = ?X' x if x ∈ codomain c for x
    using c.fst-apply-chart-TM-id X k that by (metis c comp-apply prod.collapse)
  have smooth-on-snd-chart-TM: smooth-on (codomain c) (snd ∘ ?X)
    using asm-df.diff-chartsD[OF c-in-sub-atlas, of chart-eucl] codomain-c-eq
    by (auto simp add: comp-assoc smooth-on-cong)

  show ?smooth-in-chart-TM c X
    apply (rule smooth-on-cong[OF - - X-eq])
    using smooth-on-Pair smooth-on-id smooth-on-snd-chart-TM by blast+
  qed
qed
end

```

context *smooth-vector-field-local* **begin**

definition *chart-X* ≡ λp. (c-manifold-point.tangent-chart-fun charts ∞ ψ p) (X p)

lemma *smooth-in-chart-X* [*simp*]: $\text{diff-fun} \infty (\text{charts-submanifold} (\text{domain } \psi))$
chart-X

unfolding *chart-X-def* **using** *smooth-in-chart* **by** *simp*

lemma *apply-chart-TM-chart-X*:

$\text{smooth-on} (\text{codomain } \psi) (\text{apply-chart-TM } \psi \circ (\lambda p. (p, X p)) \circ \text{inv-chart } \psi) \longleftrightarrow$
 $\text{diff-fun} \infty (\text{charts-submanifold} (\text{domain } \psi)) \text{ chart-X}$

unfolding *chart-X-def*

apply (*rule apply-chart-TM-chartX*[*of* ψ $\lambda p. (p, X p)$, *simplified*])

unfolding *section-of-TM-def in-TM-def* **apply** (*clarsimp*, *intro conjI*)

using ψ *vector-field* **by** (*blast*, *auto*)

end

5.2.2 Some theorems about smooth vector fields, locally and globally.

context *c-manifold-local* **begin**

It is often convenient to keep a stronger handle on which chart we're (locally) working in. Since the first component of the *apply-chart-TM* is just the identity, we can safely omit it for a lot of our reasoning about smoothness in a chart (see $\llbracket \text{section-of-TM-on} (\text{domain } ?\psi) ?X; ?\psi \in \text{atlas}; ?x \in \text{codomain } ?\psi \rrbracket \implies (\text{fst} \circ (\text{apply-chart-TM } ?\psi \circ ?X \circ \text{inv-chart } ?\psi)) ?x = ?x$ and $\llbracket k = \infty; \text{section-of-TM-on} (\text{domain } ?c) ?X; ?c \in \text{atlas} \rrbracket \implies \text{smooth-on} (\text{codomain } ?c) (\text{apply-chart-TM } ?c \circ ?X \circ \text{inv-chart } ?c) = \text{diff-fun} \infty (\text{charts-submanifold} (\text{domain } ?c)) (\lambda p. \text{c-manifold-point.tangent-chart-fun charts } \infty ?c p (\text{snd } (?X p)))$).

definition *vector-field-component* :: $('a \Rightarrow ((a \Rightarrow \text{real}) \Rightarrow \text{real})) \Rightarrow 'b \Rightarrow 'a \Rightarrow \text{real}$

where *vector-field-component* $X i \equiv \lambda p. (\text{c-manifold-point.component-function charts } k \psi p) (X p) i$

definition *coordinate-vector-field* :: $'b \Rightarrow ('a \Rightarrow ((a \Rightarrow \text{real}) \Rightarrow \text{real}))$

where *coordinate-vector-field* $i p \equiv \text{c-manifold-point.coordinate-vector charts } k \psi p i$

Eqn. 8.2, page 175, Lee 2012

lemma *vector-field-local-representation*:

assumes $k: k = \infty$ **and** X : *rough-vector-field* X **and** p : $p \in \text{domain } \psi$

shows $X p = (\sum_{i \in \text{Basis}. (\text{vector-field-component } X i p) *_{\mathbb{R}} (\text{coordinate-vector-field } i p))$

unfolding *vector-field-component-def* *coordinate-vector-field-def*

apply (*rule c-manifold-point.coordinate-vector-representation*)

apply *unfold-locales*

subgoal using p *rough-vector-fieldE*[*OF* X] *sub- ψ -carrier* **by** *blast*

subgoal using p *rough-vector-fieldE*[*OF* X] *in-carrier-atlasI*[*OF* ψ] **by** *blast*

by (*simp add: k*)

definition *local-coord-at* :: 'a \Rightarrow 'b \Rightarrow 'a \Rightarrow real
where *local-coord-at* q i \equiv restrict0 (domain ψ) ($\lambda y::'a. (\psi y - \psi q) \cdot i$)

lemma *local-coord-diff-fun*:
assumes k: k= ∞ **and** q: q \in domain ψ
shows *local-coord-at* q i \in sub- ψ .sub.diff-fun-space
proof –
note *local-simps*[simp] = *local-coord-at-def*
have *diff-fun* k (charts-submanifold (domain ψ)) (($\lambda y::'a. (\psi y - \psi q) \cdot i$))
apply (rule *diff-fun-compose*[unfolded o-def, of k - charts-eucl ψ])
using *diff-fun- ψ .diff-axioms* k **by** (auto intro!: *diff-fun-charts-euclI smooth-on-inner smooth-on-minus*)
from *diff-fun.diff-fun-cong*[OF this] q
have *diff-fun* k (charts-submanifold (domain ψ)) (*local-coord-at* q i) **by** simp
then show *local-coord-at* q i \in sub- ψ .sub.diff-fun-space
by auto (*metis restrict0-subset sub- ψ sub- ψ .sub.domain-atlas-subset-carrier sub- ψ .sub.restrict0-in-fun-space*)
qed

lemma *vector-apply-coord-at*:
fixes x_ψ **defines** [simp]: $x_\psi \equiv$ *local-coord-at*
assumes q: q \in domain ψ **and** p:p \in domain ψ **and** X: X \in tangent-space q **and**
k: k= ∞
shows ($d\iota^{-1}$ q) X (x_ψ p i) = ($d\iota^{-1}$ q) X (x_ψ q i)
proof –
note *local-simps*[simp] = *local-coord-at-def*
have *diff- x_ψ i'*: x_ψ q i \in sub- ψ .sub.diff-fun-space **if** q \in domain ψ **for** i q
using *local-coord-diff-fun*[OF k that] **by** simp

interpret q: c-manifold-point charts k ψ q **using** q ψ **by** *unfold-locales simp*
let ? $x_q = x_\psi$ q
have Xq: q.dRestr2 X \in q.T_pU
using *bij-betwE*[OF q.bij-betw-d ι -inv] X **by** simp
{
fix x' b **assume** x' \in domain ψ
have Dp-simp: *frechet-derivative* ((x_ψ p' i) \circ inv-chart ψ) (at (ψ x')) =
frechet-derivative (($\lambda y. y \cdot i$) (at (ψ x'))) **for** p'
proof –
have *frechet-derivative* ((x_ψ p' i) \circ inv-chart ψ) (at (ψ x')) = *frechet-derivative*
(($\lambda y. (y - \psi p') \cdot i$) (at (ψ x')))
apply (rule *frechet-derivative-transform-within-open*[OF - open-codomain[of ψ], *symmetric*])
by (*simp-all add*: $\langle x' \in$ domain $\psi \rangle$)
then show ?thesis
by (auto simp: *algebra-simps zero-fun-def*
intro!: *frechet-derivative-at*[*symmetric*] *has-derivative-diff*[**where** g'=0,
simplified] *derivative-intros*)
qed

```

have frechet-derivative ((xψ p i) ∘ inv-chart ψ) (at (ψ x')) b = frechet-derivative
((xψ q i) ∘ inv-chart ψ) (at (ψ x')) b
  by (simp only: Dp-simp)
} note deriv-eq = this
show ?thesis
  apply (rule sub-ψ.sub.derivation-eq-localI'[OF k q - - Xq, of ψ])
  using local-coord-diff-fun diff-xψi' k deriv-eq sub-ψ
  by (auto simp: p sub-ψ.sub.diff-fun-space-def)
qed

```

end

context *c-manifold* **begin**

abbreviation (input) *real-linear-on S1 S2* ≡ *linear-on S1 S2 scaleR scaleR*

— Sometimes we want to apply a vector field meaningfully to a function that is in the *c-manifold.diff-fun-space* of a submanifold (e.g. a single chart). For this to make sense, the function has to be in the correct space, and the submanifold's carrier set has to be open.

definition *vec-field-apply-fun-in-at* :: ('a vector-field) ⇒ ('a ⇒ real) ⇒ 'a set ⇒ 'a ⇒ real

```

where vec-field-apply-fun-in-at X f U q = restrict0 (tangent-space q)
  (the-inv-into
    (c-manifold.tangent-space (charts-submanifold U) k q)
    (diff.push-forward k (charts-submanifold U) charts (λx. x)))
  (X q) f

```

abbreviation *vec-field-restr* :: ('a vector-field) ⇒ 'a set ⇒ ('a vector-field)

where *vec-field-restr* X U q f ≡ restrict0 U (*vec-field-apply-fun-in-at* X f U) q

notation *vec-field-restr* (⟨-⟩) [60,60]

lemma (in *smooth-manifold*) *vec-field-restr*: (X ↯ U) p ∈ *c-manifold.tangent-space* (charts-submanifold U) ∞ p

if open U U ⊆ carrier rough-vector-field X **for** U X

proof –

interpret U: submanifold charts ∞ U

by (unfold-locales, simp add: that)

have U-simps[simp]: U.sub.carrier = U

using that **by** auto

show ?thesis

apply (cases p ∈ U)

subgoal

apply (simp add: *vec-field-apply-fun-in-at-def*)

using *bij-betwE*[OF U.bij-betw-dt-inv] that rough-vector-fieldE(1) **by** auto

by (simp add: U.sub.diff-fun-space.linear-zero U.sub.tangent-spaceI extensional0-def)

qed

lemma *vec-field-apply-fun-alt'*:

assumes *open U q ∈ U f ∈ c-manifold.diff-fun-space (charts-submanifold U) k rough-vector-field X*

shows *vec-field-apply-fun-in-at X f U q = (the-inv-into (c-manifold.tangent-space (charts-submanifold U) k q) (diff.push-forward k (charts-submanifold U) charts (λx. x))) (X q) f*

using *rough-vector-fieldE(1)[OF assms(4)] by (auto simp: vec-field-apply-fun-in-at-def assms(1-3))*

lemma *vec-field-apply-fun-alt*:

assumes *open U q ∈ U f ∈ c-manifold.diff-fun-space (charts-submanifold U) k rough-vector-field X*

shows *vec-field-restr X U q f = (the-inv-into (c-manifold.tangent-space (charts-submanifold U) k q) (diff.push-forward k (charts-submanifold U) charts (λx. x))) (X q) f*

using *rough-vector-fieldE(1)[OF assms(4)] by (auto simp: vec-field-apply-fun-in-at-def assms(1-3))*

lemma (*in submanifold*) *vec-field-apply-fun-sub*:

assumes *q ∈ carrier q ∈ S f ∈ sub.diff-fun-space rough-vector-field X*

shows *vec-field-apply-fun-in-at X f (S ∩ carrier) q = (the-inv-into (sub.tangent-space q) inclusion.push-forward) (X q) f*

using *assms charts-submanifold-Int-carrier sub.open-carrier vec-field-apply-fun-alt by auto*

lemma *vec-field-apply-fun-in-open[simp]*: *vec-field-apply-fun-in-at X f' U p = X p f*

if *U: p ∈ U open U U ⊆ carrier*

and *f: f ∈ diff-fun-space f' ∈ c-manifold.diff-fun-space (charts-submanifold U) k ∀ x ∈ U. f x = f' x*

and *X: rough-vector-field X*

proof –

interpret *U: submanifold charts k U using U(2) by unfold-locales*

show *?thesis*

using *U.vec-field-apply-fun-sub[OF subsetD[OF U(3,1)] U(1) f(2) X] U.vector-apply-sub-eq-localI(2)*

using *rough-vector-fieldE(1)[OF X] that(1,3-6) by (auto simp: Int-absorb2[OF U(3)] U.open-submanifold)*

qed

lemma *open-imp-submanifold*: *submanifold charts k S if open S*

using *that by unfold-locales*

lemmas *charts-submanifold = submanifold.charts-submanifold[OF open-imp-submanifold]*

lemma *charts-submanifold-Int*:

manifold.charts-submanifold (charts-submanifold U) N = charts-submanifold (N ∩ U)

if *open N open U*

using *restrict-chart-restrict-chart*[*OF that*] **by** (*auto simp add: image-def manifold.charts-submanifold-def*)

lemma *vec-field-apply-fun-in-restrict0*[*simp*]:

vec-field-restr $X\ U\ p\ f = \text{vec-field-restr}\ X\ N\ p\ (\text{restrict0}\ N\ f)$

if U : *open* $U\ U \subseteq \text{carrier}$ **and** N : $p \in N\ N \subseteq U$ *open* N

and f : $f \in \text{c-manifold.diff-fun-space}\ (\text{charts-submanifold}\ U)\ k$

and X : *rough-vector-field* X

proof –

let $?f = \text{restrict0}\ N\ f$

have $f\text{-diff-N}$: *diff-fun* $k\ (\text{charts-submanifold}\ N)\ f$

using *diff-fun.diff-fun-submanifold*[*OF c-manifold.diff-fun-spaceD* [*OF charts-submanifold* [*OF U*(1)]]], *OF f N*(3)]

by (*simp only: charts-submanifold-Int* [*OF N*(3) *U*(1)] *Int-absorb2* [*OF N*(2)])

have f' : $?f \in \text{c-manifold.diff-fun-space}\ (\text{charts-submanifold}\ N)\ k$

unfolding *c-manifold.diff-fun-space-def* [*OF charts-submanifold* [*OF N*(3)]] **apply**

(*safe, rule diff-fun.diff-fun-cong*)

using *f-diff-N submanifold.carrier-submanifold* [*OF open-imp-submanifold* [*OF N*(3)]]

by (*auto simp: Int-absorb2* [*OF subset-trans, OF N*(2) *U*(2)])

have p : $p \in U\ p \in \text{carrier}$ **using** $U\ N$ **by** *auto*

have $N\text{-carrier}$ [*simp*]: *manifold.carrier* (*charts-submanifold* N) = N

using *submanifold.carrier-submanifold open-imp-submanifold N*(3) *Int-absorb2 N*(2) *U*(2)

by (*metis subset-trans*)

obtain N' **where** N' : $p \in N'$ *open* N' *compact* (*closure* N') *closure* $N' \subseteq N$

using *manifold.precompact-neighborhoodE* [*of p charts-submanifold N, simplified, OF N*(1)] **by** *blast*

from *submanifold.extension-lemma-submanifoldE* [*OF open-imp-submanifold* [*OF N*(3)]] *f-diff-N closed-closure*] *this*(4)

obtain g **where** g : *diff-fun* $k\ \text{charts}\ g \wedge x. x \in \text{closure}\ N' \implies g\ x = f\ x$

csupport-on carrier g $\cap \text{carrier} \subseteq N$

by *auto*

let $?g = \text{restrict0}\ \text{carrier}\ g$

have $\text{diff-g}'$: *diff-fun* $k\ \text{charts}\ ?g\ ?g \in \text{diff-fun-space}$

subgoal **by** (*rule diff-fun.diff-fun-cong* [*OF g*(1)]) *simp*

subgoal **unfolding** *diff-fun-space-def* **using** $\langle \text{diff-fun}\ k\ \text{charts}\ ?g \rangle$ **by** *simp*

done

note [*simp*] = *charts-submanifold-Int* [*OF N*(3) *U*(1)] *Int-absorb2* [*OF N*(2)] *rough-vector-fieldE*(1) [*OF X*]

vec-field-apply-fun-alt [*OF N*(3,1) f'] *vec-field-apply-fun-alt* [*OF U*(1) p (1) f]

note $Xp\text{-eq-localI}$ = *submanifold.vector-apply-sub-eq-localI*(2)

[*OF open-imp-submanifold N*'(1) - N' (2)]

subset-trans [*OF closure-subset, OF subset-trans* [*OF N*'(4)]]

diff-g'(2) - - *rough-vector-fieldE*(1) [*OF X*]]

have $f\text{-eq}$: $\text{restrict0 carrier } g \ x = f \ x \ \text{restrict0 carrier } g \ x = \text{restrict0 } N \ f \ x$
if $x \in N'$ **for** x
proof –
have $x \in \text{carrier } x \in N$
using $\text{that } N'(4) \ N(2) \ U(2)$ **by** auto
thus $\text{restrict0 carrier } g \ x = f \ x \ \text{restrict0 carrier } g \ x = \text{restrict0 } N \ f \ x$
using $\text{that } g(2)$ **by** auto
qed

show $?thesis$
using $Xp\text{-eq-localI}[OF \ N(3) \ \text{subset-trans}[OF \ N(2)], \ OF \ U(2) - f' \ f\text{-eq}(2)]$
 $Xp\text{-eq-localI}[OF \ U \ N(2) \ f \ f\text{-eq}(1)]$
by $(\text{simp add: } X \ f' \ \text{that}(3) \ \text{that}(5) \ \text{vec-field-apply-fun-alt'})$
qed

lemma (**in** submanifold) $\text{vec-field-apply-fun-in-open}[\text{simp}]$:
 $\text{vec-field-restr } X \ S \ p \ f' = X \ p \ f$
if $S: S \subseteq \text{carrier}$
and $N: \text{open } N \ N \subseteq S \ p \in N$
and $f: f \in \text{diff-fun-space } f' \in \text{sub.diff-fun-space } \forall x \in N. f \ x = f' \ x$
and $X: \text{rough-vector-field } X$
using $\text{vector-apply-sub-eq-localI}(2)[OF \ N(3) \ S \ N(1,2) \ f(1,2)] \ \text{that}(3,4,6,7,8)$
by $(\text{auto simp: vec-field-apply-fun-alt' rough-vector-fieldE}(1) \ \text{open-submanifold})$

lemma (**in** smooth-manifold) $\text{vec-field-apply-fun-in-restrict0}'$:
 $\text{restrict0 } U \ (Xf) = X \upharpoonright U'' \ (\text{restrict0 } U \ f)$
if $U: \text{open } U \ U \subseteq \text{carrier}$ **and** $f: f \in \text{diff-fun-space}$ **and** $X: \text{rough-vector-field } X$
for $U \ X \ f$
proof
fix p
interpret $U: \text{submanifold charts } \infty \ U$
by $(\text{unfold-locales, simp add: } U)$
have $U\text{-simps}[\text{simp}]: U.\text{sub.carrier} = U$
using U **by** auto

show $\text{restrict0 } U \ (Xf) \ p = X \upharpoonright U \ p \ (\text{restrict0 } U \ f)$ (**is** $\langle ?LHS = ?RHS \rangle$)
by $(\text{metis } (\text{mono-tags, lifting}) \ U.\text{open-submanifold } X \ U(2) \ \text{charts-submanifold-carrier}$
 $\text{diff.defined } \text{diff-id } f \ \text{image-ident } \text{open-carrier } \text{open-imp-submanifold } \text{re-$
 strict0-def
 $\text{submanifold.vec-field-apply-fun-in-open } \text{vec-field-apply-fun-in-restrict0})$
qed

lemma (**in** submanifold) $\text{vec-field-apply-fun-in-open}'[\text{simp}]$:
 $\text{vec-field-restr } X \ S \ p \ f' = X \ p \ f$
if $S: p \in S \ S \subseteq \text{carrier}$
and $f: f \in \text{diff-fun-space } f' \in \text{sub.diff-fun-space } \forall x \in S. f \ x = f' \ x$

```

and  $X$ : rough-vector-field  $X$ 
using vec-field-apply-fun-in-open[ $OF\ S(2)\ open-submanifold - S(1)\ f\ X$ ] by simp

lemma (in  $c$ -manifold) vec-field-apply-fun-in-chart[simp]:
  vec-field-apply-fun-in-at  $X\ f\ (domain\ c)\ p = X\ p\ f$ 
  if  $p$ :  $p \in domain\ c$  and  $c$ :  $c \in atlas$ 
    and  $f$ :  $f \in diff-fun-space\ f \in c-manifold.diff-fun-space\ (charts-submanifold\ (domain\ c))\ k$ 
    and  $X$ : rough-vector-field  $X$ 
    using vec-field-apply-fun-in-open that by blast

end

```

```

context  $c$ -manifold-local begin

```

```

lemma vec-field-apply-fun-eq-component:
  fixes  $x_\psi$  defines [simp]:  $x_\psi \equiv local-coord-at$ 
  assumes  $q$ :  $q \in domain\ \psi$  and  $p$ :  $p \in domain\ \psi$  and  $X$ : rough-vector-field  $X$  and
   $k$ :  $k = \infty$ 
  shows vec-field-apply-fun-in-at  $X\ (x_\psi\ q\ i)\ (domain\ \psi)\ q = vector-field-component\ X\ i\ q$ 
  proof –
  note [simp] = local-coord-at-def sub- $\psi$ .sub.diff-fun-space-def vector-field-component-def
  interpret  $q$ :  $c$ -manifold-point charts  $k\ \psi\ q$  using  $q\ \psi$  by unfold-locales simp
  let  $?x_q = x_\psi\ q$ 
  have  $Xq$ :  $X\ q \in q.T_pM\ q.dRestr2\ (X\ q) \in q.T_pU$ 
    subgoal using rough-vector-fieldE[ $OF\ X$ ]  $q\ \psi$  by blast
    using bij-betwE[ $OF\ q.bij-betw-dt-inv$ ]  $\langle X\ q \in q.T_pM \rangle$  by simp
  note  $1 = vector-apply-coord-at[OF\ q\ p\ Xq(1)\ k]$ 
  have  $2$ :  $q.dRestr2\ (X\ q)\ (local-coord-at\ q\ i) = vector-field-component\ X\ i\ q$ 
    using  $q.component-function-apply-in-T_pM[OF\ Xq(1)]$  by simp
  show  $?thesis$ 
    apply (simp only:  $2[symmetric]\ 1[symmetric]\ restrict0-apply-in[OF\ Xq(1)]$ )
    using vec-field-apply-fun-alt'[ $OF\ open-domain\ q$ ] local-coord-diff-fun[ $OF\ k\ q$ ]  $X$ 
     $x_\psi$ -def
    by blast
qed

```

Prop 8.1, page 175, Lee 2012. The main difference is that our vector field X here is only a map $M \rightarrow snd^i TM$, not a section $M \rightarrow TM$ properly speaking. See also $\llbracket k = \infty; section-of-TM-on\ (domain\ ?c)\ ?X; ?c \in atlas \rrbracket \implies smooth-on\ (codomain\ ?c)\ (apply-chart-TM\ ?c \circ ?X \circ inv-chart\ ?c) = diff-fun\ \infty\ (charts-submanifold\ (domain\ ?c))\ (\lambda p. c-manifold-point.tangent-chart-fun\ charts\ \infty\ ?c\ p\ (snd\ (?X\ p)))$.

```

lemma vector-field-smooth-local-iff:
  assumes  $k$ :  $k = \infty$  and  $X$ :  $\forall p \in domain\ \psi. X\ p \in tangent-space\ p$ 

```

shows *smooth-vector-field-local charts* $\psi X \longleftrightarrow (\forall i \in \text{Basis}. \text{diff-fun-on } (\text{domain } \psi))$ (*vector-field-component* $X i$)

(**is** $\langle ?\text{smooth-vf } X \longleftrightarrow (\forall i \in \text{Basis}. ?\text{diff-component } X i) \rangle$)

proof –

A bit of house-keeping. Maybe having both *vector-field-component* and *c-manifold-point.tangent-chart-fun* is redundant, or maybe the second statement below could be a simp rule (probably in the opposite direction).

have *cpt1: c-manifold-point charts* $k \psi a$ **if** $a \in \text{domain } \psi$ **for** a

apply *unfold-locale* **by** (*simp-all add: sub- ψ that*)

have *vfc-tcf: $(\sum i \in \text{Basis}. \text{vector-field-component } X i p *_R i) = \text{c-manifold-point.tangent-chart-fun charts } \infty \psi p (X p)$*

if $p \in \text{domain } \psi$ **for** p

using *c-manifold-point.tangent-chart-fun-def[*of charts* $k \psi$] vector-field-component-def cpt1 k that* **by** *auto*

show *?thesis*

proof

assume *asm: ?smooth-vf X*

then interpret *smooth- X -local: smooth-vector-field-local charts* ψX

unfolding *smooth-vector-field-local-def* .

Notice we don't even need our Euclidean representation theorem $\llbracket k = \infty; \text{rough-vector-field } ?X; ?p \in \text{domain } \psi \rrbracket \implies ?X ?p = (\sum i \in \text{Basis}. \text{vector-field-component } ?X i ?p *_R \text{coordinate-vector-field } i ?p)$. The crux is that we already know differentiability works well with components: *diff-fun k charts $(\lambda x. \sum i \in \text{Basis}. ?f i x *_R i) = (\forall i \in \text{Basis}. \text{diff-fun } k \text{ charts } (?f i))$* .

have $\forall i \in \text{Basis}. \text{diff-fun } \infty (\text{charts-submanifold } (\text{domain } \psi)) (\text{vector-field-component } X i)$

apply (*subst smooth- X -local.sub- ψ .sub.diff-fun-components-iff[*of vector-field-component* X , *symmetric*]*)

using *smooth- X -local.smooth-in-chart- X [*unfolded smooth- X -local.chart- X -def*]*

by (*auto intro: diff-fun.diff-fun-cong[OF - vfc-tcf[*symmetric*]]*)

then show $\forall i \in \text{Basis}. ?\text{diff-component } X i$

using *diff-fun-onI[*of domain* ψ *domain* ψ *ff* **for** f] domain-atlas-subset-carrier k by auto*

next

assume *asm: $\forall i \in \text{Basis}. ?\text{diff-component } X i$*

interpret *local- ψ : c-manifold-local charts* $\infty \psi$ **using** *c-manifold-local-axioms k by auto*

have *2: diff-fun $\infty (\text{charts-submanifold } (\text{domain } \psi)) (\lambda p. \text{c-manifold-point.tangent-chart-fun charts } \infty \psi p (X p))$*

apply (*rule diff-fun.diff-fun-cong[OF - vfc-tcf]*)

using *sub- ψ .sub.diff-fun-components-iff[*of vector-field-component* X] k asm diff-fun-on-open*

by *auto*


```

have 1: smooth-on (codomain  $\psi$ ) (( $\lambda p$ . c-manifold-point.tangent-chart-fun charts
 $\circ \psi p (X p)$ )  $\circ$  inv-chart  $\psi$ )
  if  $x \in$  domain  $\psi$  for  $x$ 
  apply (rule diff-fun.diff-fun-between-chartsD[of - charts-submanifold (domain
 $\psi$ )])
  using 2 that by (auto simp: local- $\psi$ .sub- $\psi$ )

show ?smooth-vf  $X$ 
  apply unfold-locales
  using 1  $X$  by (auto intro!: bexI[of -  $\psi$ ] bexI[of - chart-eucl] simp: local- $\psi$ .sub- $\psi$ 
id-comp[unfolded id-def])
qed
qed

```

end

```

lemma (in smooth-vector-field-local) diff-component':
  fixes  $i :: 'b$ 
  assumes  $i \in$  Basis
  shows diff-fun-on (domain  $\psi$ ) (vector-field-component  $X i$ )
  using vector-field-smooth-local-iff[OF - vector-field] smooth-vector-field-local-axioms
  assms by auto

```

context *smooth-manifold* **begin**

Prop. 8.8 in Lee 2012.

Do we want extensional0 vector fields? It would make the usual simplification for writing addition and scaling by real numbers. So \mathfrak{X} could be a vector space under (+) and *scaleR*? Maybe a double problem: * ~~0 is ill-defined when $'a$ is not of the sort *zero*.~~ * Also I think the function 0 always assigns zero, i.e. for a pair it returns the constant $(0,0)$. We would want the zero vector field to be $p \mapsto (p, 0)$ instead.

We will need to use locales anyway if we also want to talk about \mathfrak{X} as a module over *diff-fun-space*, since that is a set already. - Actually, probably not true, because *extensional0* works out quite neatly.

A predicate analogous to *smooth-vector-field-local*, but for the entire manifold.

```

definition smooth-vector-field :: ' $a$  vector-field  $\Rightarrow$  bool
  where smooth-vector-field  $X \equiv$  rough-vector-field  $X \wedge$  smooth-from-M-to-TM
( $\lambda p$ .  $(p, X p)$ )

```

```

lemma smooth-vector-field-alt:
  smooth-vector-field  $X \equiv$  ( $\lambda p$ .  $(p, X p)$ )  $\in$   $\mathfrak{X} \wedge$  extensional0 carrier  $X$ 

```

by (*auto simp: smooth-vector-field-def smooth-section-of-TM-def section-of-TM-def rough-vector-field-def in-TM-def, linarith*)

— For displaying.

lemma *smooth-vector-field* $X \equiv (\forall p \in \text{carrier}. X p \in \text{tangent-space } p) \wedge$
 $\text{smooth-from-M-to-TM } (\lambda p. (p, X p)) \wedge$
 $\text{extensional0 carrier } X$

by (*auto simp: smooth-vector-field-def smooth-section-of-TM-def section-of-TM-def rough-vector-field-def in-TM-def, linarith*)

lemma *smooth-vector-fieldE* [*elim*]:

assumes *smooth-vector-field* X

shows $\bigwedge p. X p \in \text{tangent-space } p \text{ extensional0 carrier } X \text{ rough-vector-field } X$
 $\text{smooth-from-M-to-TM } (\lambda p. (p, X p))$

using *rough-vector-fieldE* *assms* **unfolding** *smooth-vector-field-def* **by** *auto*

lemma *smooth-vector-field-imp-local*:

assumes *smooth-vector-field* $X \psi \in \text{atlas}$

shows *smooth-vector-field-local charts* ψX

proof —

interpret ψ : *c-manifold-local charts* $\infty \psi$ **using** *assms(2)* **by** *unfold-locales*

have 1: *section-of-TM* $(\lambda p. (p, X p))$

using *assms(1,2)* *smooth-section-of-TM-def* *smooth-vector-field-alt* **by** *blast*

have 2: *smooth-from-M-to-TM* $(\lambda p. (p, X p))$

using *assms(1)* *smooth-vector-field-def* **by** *auto*

have *vec-field-X*: $\forall p \in \text{domain } \psi. X p \in \text{tangent-space } p$

using *assms(1)* **by** (*auto simp: smooth-vector-field-def*)

moreover **have** *diff-fun-X*: *diff-fun* $\infty (\text{charts-submanifold } (\text{domain } \psi)) (\lambda p.$
 $\text{c-manifold-point.tangent-chart-fun charts } \infty \psi p (X p))$

using *apply-chart-TM-chartX* *diff-from-M-to-TM-chartsD[OF 2 1, of ψ]* *section-of-TM-subset[OF 1]*

using $\psi.\psi$ **by** (*simp add: domain-atlas-subset-carrier*)

ultimately show *?thesis*

using $\psi.\text{c-manifold-local-axioms}$ **by** (*auto simp: smooth-vector-field-local-def smooth-vector-field-local-axioms-def*)

qed

lemma *smooth-vector-field-imp-local'*:

fixes $X \psi X_\psi$ **defines** $X_\psi \equiv \text{restrict0 } (\text{domain } \psi) X$

assumes *smooth-vector-field* $X \psi \in \text{atlas}$

shows *smooth-vector-field-local charts* ψX_ψ

unfolding *smooth-vector-field-local-def* *smooth-vector-field-local-axioms-def* *assms(1)*

using *smooth-vector-field-imp-local[OF assms(2-)]* **apply** *safe*

subgoal using *smooth-vector-field-local.axioms(1)* **by** *blast*

subgoal using *rough-vector-fieldE(1)* *smooth-vector-field-local.rough-vector-field*
by *blast*

apply (*rule* *diff-fun.diff-fun-cong[of - - ($\lambda p. \text{c-manifold-point.tangent-chart-fun charts } \infty \psi p (X p))]$)*

subgoal by (*simp add: assms(2,3)* *smooth-vector-field-imp-local smooth-vector-field-local.smooth-in-chart*)

subgoal by (*metis IntD2 inf-sup-aci(1) open-domain open-imp-submanifold restrict0-apply-in submanifold.carrier-submanifold*)

done

lemma *smooth-vector-field-if-local*:

assumes $\forall p \in \text{carrier}. \exists c \in \text{atlas}. p \in \text{domain } c \wedge \text{smooth-vector-field-local charts } c \ X \ \text{extensional0 carrier } X$

shows *smooth-vector-field X*

proof (*unfold smooth-vector-field-def diff-from-M-to-TM-def, intro conjI allI impI*)

show *vec-field-X: rough-vector-field X*

by (*metis assms(1,2) rough-vector-field-def smooth-vector-field-local.vector-field*)

fix *p* **assume** *p* $\in \text{carrier}$

then obtain *c* **where** *c*: $c \in \text{atlas}$ **and** *p*: $p \in \text{domain } c$ **and** *X*: *smooth-vector-field-local charts c X*

using *assms(1)* **by** *blast*

interpret *X*: *smooth-vector-field-local charts c X* **using** *X* .

have *im-domain*: $(\lambda p. (p, X p)) \text{ ' domain } c \subseteq \text{domain-TM } c$

using *rough-vector-fieldE[OF vec-field-X] in-carrier-atlasI[OF c]* **by** (*auto simp: domain-TM-def*)

have *smooth-X*: *smooth-on (codomain c) (apply-chart-TM c \circ ($\lambda p. (p, X p)$) \circ inv-chart c)*

using *X.apply-chart-TM-chart-X X.smooth-in-chart-X* **by** *blast*

show $\exists c \in \text{atlas}. p \in \text{domain } c \wedge (\lambda p. (p, X p)) \text{ ' domain } c \subseteq \text{domain-TM } c \wedge \text{smooth-on (codomain } c) (\text{apply-chart-TM } c \circ (\lambda p. (p, X p))) \circ \text{inv-chart } c$

using *c p im-domain smooth-X* **by** *blast*

qed

lemma *smooth-vector-field-iff-local*:

assumes *extensional0 carrier X*

shows $(\forall c \in \text{atlas}. \text{smooth-vector-field-local charts } c \ X) \longleftrightarrow \text{smooth-vector-field } X$

using *smooth-vector-field-imp-local smooth-vector-field-if-local* **by** (*meson assms atlasE*)

— For display mostly.

lemma (*in smooth-manifold*) *smooth-vector-field-local*:

assumes $c \in \text{atlas} \ \forall p \in \text{domain } c. X \ p \in \text{tangent-space } p$

shows *smooth-vector-field-local charts c X* \longleftrightarrow

smooth-on (codomain c) (apply-chart-TM c \circ ($\lambda p. (p, X p)$) \circ inv-chart c)

proof –

```

interpret c: submanifold charts ∞ domain c
  using open-domain by unfold-locales simp
let ?c1 = λx. c (inv-chart c x)
let ?c2 = λx. c-manifold-point.tangent-chart-fun charts ∞ c (inv-chart c x) (X (inv-chart c x))
{
  fix x assume smoothTM: smooth-on (codomain c) (λx. (?c1 x, ?c2 x)) and x:
x ∈ c.sub.carrier
  have loc-simp: snd ∘ (λx. (?c1 x, ?c2 x)) = ?c2 by auto
  have x ∈ domain c ∧ c ∈ c.sub.atlas ∧ smooth-on (codomain c) ?c2
  using x apply (simp add: assms(1) atlas-is-atlas c.sub.maximal-atlas c.submanifold-atlasE domain-atlas-subset-carrier)
  apply (subst loc-simp[symmetric, unfolded o-def])
  apply (rule smooth-on-snd[of ∞, OF - open-codomain[of c]])
  using smoothTM .
}
thus ?thesis
  unfolding smooth-vector-field-local-def smooth-vector-field-local-axioms-def
  apply (intro iffI)
  using smooth-vector-field-local.apply-chart-TM-chart-X smooth-vector-field-local.intro smooth-vector-field-local.smooth-in-chart-X smooth-vector-field-local-axioms-def apply blast
  apply (simp add: assms c-manifold-axioms c-manifold-local.intro c-manifold-local-axioms.intro)
  apply (rule c-manifold.diff-funI[OF charts-submanifold, OF open-domain[of c]])
  unfolding apply-chart-TM-def apply (simp add: o-def)
  apply (rule bexI[of - c c-manifold.atlas (charts-submanifold (domain c)) ∞])
  by blast+
qed

```

lemma (**in** *c-manifold*) *diff-fun-deriv-chart'*:

```

fixes i::'b
assumes c:c∈atlas and f:diff-fun-on (domain c) f and k: k>0
shows diff-fun (k-1) (charts-submanifold (domain c)) (λx. frechet-derivative (f
  ∘ inv-chart c) (at (c x) i))
proof –
  have local-simps [simp]: k - 1 + 1 = k
  using k by (metis add commute add-diff-assoc-enat add-diff-cancel-enat ileI1 infinity-ne-i1 one-eSuc)
  interpret c1: c-manifold-local charts k-1 c
  apply unfold-locales
  apply (metis c-manifold-def c-manifold-order-le le-iff-add local-simps)
  by (metis c in-atlas-order-le le-iff-add local-simps)
  interpret f': diff-fun k charts-submanifold (domain c) f
  using diff-fun-on-open[of domain c f] f by simp
  { fix x and j::'b assume x: x∈domain c x∈carrier
    have (k-1)–smooth-on (codomain c) (λy. frechet-derivative (f ∘ (inv-chart c))
```

(at y) j)
apply (rule derivative-is-smooth'[of - codomain c], simp)
apply (rule f'.diff-fun-between-chartsD)
using c c-manifold-local.sub-ψ c-manifold-point c-manifold-point.axioms(1) k
x(1) **by** blast+
then have (k-1)-smooth-on (codomain c) (λx. frechet-derivative (λx. f
(inv-chart c x)) (at (c (inv-chart c x)))) j)
by (auto intro: smooth-on-cong simp: o-def) }
then show ?thesis
by (auto intro!: c1.sub-ψ.sub.diff-funI bexI[of - c] simp: o-def c1.sub-ψ)
qed

lemma diff-fun-deriv-chart:
fixes i::'b
assumes c:c∈atlas **and** f:diff-fun-on (domain c) f
shows diff-fun ∞ (charts-submanifold (domain c)) (λx. frechet-derivative (f ∘
inv-chart c) (at (c x)) i)
using diff-fun-deriv-chart'[OF assms] **by** auto

lemma (in c-manifolds) diff-localI2: diff k charts1 charts2 f
if ∀ x∈src.carrier. (∃ U. diff k (src.charts-submanifold U) charts2 f ∧ open U ∧
x ∈ U)
using diff-localI that **by** metis

5.3 Smooth vector fields as maps $C^\infty(M) \rightarrow C^\infty(M)$.

Proposition 8.14 in Lee 2012.

lemma vector-field-smooth-iff:
assumes X: rough-vector-field X
shows smooth-vector-field X \longleftrightarrow (∀ f∈diff-fun-space. (X'' f) ∈ diff-fun-space)
(is ⟨?LHS \longleftrightarrow ?RHS1⟩)
and smooth-vector-field X \longleftrightarrow (∀ U f. open U ∧ U ⊆ carrier ∧ f∈(c-manifold.diff-fun-space
(charts-submanifold U) ∞) \longrightarrow
diff-fun ∞ (charts-submanifold U))
(vec-field-apply-fun-in-at X f U))
(is ⟨?LHS \longleftrightarrow ?RHS2⟩)

proof –

Prove a chain of implications smooth-vector-field X \longrightarrow (∀ f∈diff-fun-space.
(λp. X p f) ∈ diff-fun-space) \longrightarrow (∀ U f. open U ∧ U ⊆ carrier ∧ f ∈
c-manifold.diff-fun-space (charts-submanifold U) ∞ \longrightarrow diff-fun ∞ (charts-submanifold
U) (vec-field-apply-fun-in-at X f U)) \longrightarrow smooth-vector-field X to conclude
both equivalences, following Lee 2012, pp. 180–181.

have ?RHS1 **if** smooth-X: ?LHS

proof

fix f **assume** f: f ∈ diff-fun-space

{ **fix** p **assume** p: p∈carrier

obtain c **where** $c: p \in \text{domain } c \ c \in \text{atlas}$ **using** $\text{atlasE } p$ **by** blast
interpret $p: c\text{-manifold-point charts} \infty c \ p$ **by** ($\text{simp add: } p \ c\text{-manifold-point}$
 c)
{ **fix** x **assume** $x: x \in \text{domain } c$
interpret $x: c\text{-manifold-point charts} \infty c \ x$ **by** ($\text{simp add: } x \ c\text{-manifold-point}$)
have $Xx\text{-1: } X \ x \ f = (\sum_{i \in \text{Basis.}} p.\text{vector-field-component } X \ i \ x \ *_{\mathbb{R}}$
 $p.\text{coordinate-vector-field } i \ x) \ f$
by ($\text{simp only: } p.\text{vector-field-local-representation}[OF - X \ x]$)
then have $Xx\text{-2: } X \ x \ f = (\sum_{i \in \text{Basis.}} p.\text{vector-field-component } X \ i \ x \ *_{\mathbb{R}}$
 $(\text{frechet-derivative } (f \circ \text{inv-chart } c) \text{ (at } (c \ x)) \ i))$
using $x.\text{coordinate-vector-apply-in}[OF - f]$ **by** ($\text{simp add: } \text{sum-apply}$
 $p.\text{coordinate-vector-field-def}$)
}
then have $Xx\text{f: } X \ x \ f = (\sum_{i \in \text{Basis.}} p.\text{vector-field-component } X \ i \ x \ *_{\mathbb{R}}$
 $\text{frechet-derivative } (f \circ \text{inv-chart } c) \text{ (at } (c \ x)) \ i)$
if $x \in p.\text{sub-}\psi.\text{sub.carrier}$ **for** x **using** **that** **by** simp
have $\text{diff-components-}X: (\forall i \in \text{Basis. } \text{diff-fun-on } (\text{domain } c) \ (p.\text{vector-field-component}$
 $X \ i))$
using $p.\text{vector-field-smooth-local-iff } \text{rough-vector-field-subset}[OF \ X] \ \text{smooth-}X$
 $\text{domain-atlas-subset-carrier } p.\psi \ \text{smooth-vector-field-imp-local}$
by ($\text{meson } \text{smooth-vector-field-local.diff-component}'$)
have $\text{diff-fun-on } (\text{domain } c) \ f$
using $\text{diff-fun.diff-fun-submanifold}[OF \ \text{diff-fun-space}D[OF \ f]]$
by ($\text{simp add: } \text{diff-fun-on-open } \text{domain-atlas-subset-carrier}$)
note $\text{diff-fun-deriv-chart-f} = \text{diff-fun-deriv-chart}[OF \ c(2) \ \text{this}]$
have $\text{diff-}X\text{f: } \text{diff-fun} \infty (\text{charts-submanifold } (\text{domain } c)) \ (X\text{f})$
apply ($\text{rule } \text{diff-fun.diff-fun-cong}[OF - X\text{f}[\text{symmetric}]]$)
apply ($\text{rule } p.\text{sub-}\psi.\text{sub.diff-fun-sum}$)
apply ($\text{rule } p.\text{sub-}\psi.\text{sub.diff-fun-scale}\mathbb{R}$)
using $\text{diff-components-}X \ \text{diff-fun-deriv-chart-f}$ **by** ($\text{simp-all add: } \text{diff-fun-on-open}$)
have $\text{smooth-on } (\text{codomain } c) \ ((X\text{f}) \circ \text{inv-chart } c)$
using $\text{diff-}X\text{f}$ **apply** ($\text{rule } \text{diff-fun.diff-fun-between-charts}D$)
using $p.\text{sub-}\psi \ c(1)$ **by** (simp, blast)
hence $\exists c1 \in \text{atlas. } p \in \text{domain } c1 \wedge \infty\text{-smooth-on } (\text{codomain } c1) \ ((X\text{f}) \circ$
 $\text{inv-chart } c1)$
using c **by** blast }
moreover have $\text{extensional0 carrier } (X'' \ f)$
using $\text{rough-vector-fieldE}(2)[OF \ X]$ **by** ($\text{simp add: } \text{extensional0-def}$)
ultimately show $(X'' \ f) \in \text{diff-fun-space}$
unfolding $\text{diff-fun-space-def}$ **by** ($\text{auto intro: } \text{diff-funI}$)
qed

moreover have $?RHS2$ **if** $?RHS1$
proof (safe)
fix $U \ f$
assume $U: \text{open } U \ U \subseteq \text{carrier}$
and $f: f \in c\text{-manifold.diff-fun-space } (\text{charts-submanifold } U) \infty$
interpret $U: \text{submanifold charts} \infty U$ **using** $U(1)$ **by** $\text{unfold-locales simp}$

```

show  $\text{diff-fun} \infty (\text{charts-submanifold } U) (\text{vec-field-apply-fun-in-at } X f U)$ 
proof (intro  $U.\text{sub.manifold-eucl.diff-localI2 ballI}$ )
  fix  $x$  assume  $x: x \in U.\text{sub.carrier}$ 
  from  $U.\text{sub.precompact-neighborhoodE}[OF \text{ this}]$ 
  obtain  $C$  where  $C: x \in C$  open  $C$  compact (closure  $C$ )  $\text{closure } C \subseteq$ 
 $U.\text{sub.carrier}$  .
  from  $U.\text{extension-lemma-submanifoldE}[OF U.\text{sub.diff-fun-spaceD}[OF f] \text{ closed-closure}$ 
 $C(4)]$ 
  obtain  $f'$  where  $f': \text{diff-fun} \infty \text{charts } f' (\wedge x. x \in \text{closure } C \implies f' x = f x)$ 
  c-support-on carrier  $f' \cap \text{carrier} \subseteq U.\text{sub.carrier}$  by blast

  let  $?f' = \text{restrict0 } U f'$ 
  have  $1: ?f' \in \text{diff-fun-space}$ 
  unfolding  $\text{diff-fun-space-def}$  apply safe
  subgoal
  apply (rule  $\text{diff-fun.diff-fun-cong}[OF f'(1)]$ )
  using  $f'(2,3)$  by (metis (no-types)  $\text{IntE IntI } U.\text{carrier-submanifold}$ 
not-in-csupportD restrict0-def subset-iff)
  subgoal
  using  $U(2)$  extensional0-subset by blast
  done

  show  $\exists C. \text{diff} \infty (U.\text{sub.charts-submanifold } C) \text{charts-eucl} (\text{vec-field-apply-fun-in-at}$ 
 $X f U) \wedge$ 
   $\text{open } C \wedge x \in C$ 
  proof (intro  $\text{exI conjI}$ )
  have  $\text{carrier-SubC}$  [simp]:  $\text{manifold.carrier} (U.\text{sub.charts-submanifold } C)$ 
 $= C$ 
  by (metis  $C(2,4)$  Int-absorb2  $U.\text{sub.open-imp-submanifold closure-subset}$ 
 $\text{submanifold.carrier-submanifold subset-trans}$ )
  have  $\text{diff-fun} \infty (U.\text{sub.charts-submanifold } C) (\text{vec-field-apply-fun-in-at } X f$ 
 $U)$ 
  proof (rule  $\text{diff-fun.diff-fun-cong}[\text{where } f = X ?f']$ )
  have  $\text{diff-fun} \infty \text{charts} (\lambda p. X p ?f')$  using  $1$  that by (simp add:
 $\text{diff-fun-spaceD}$ )
  from  $\text{diff-fun.diff-fun-submanifold}[OF \text{ this}]$ 
  show  $\text{diff-fun} \infty (U.\text{sub.charts-submanifold } C) (\lambda p. X p ?f')$ 
  by (simp add:  $C(2)$   $U.\text{open-submanifold charts-submanifold-Int open-Int}$ )
  show  $X x (\text{restrict0 } U f') = \text{vec-field-apply-fun-in-at } X f U x$ 
  if  $x \in \text{manifold.carrier} (U.\text{sub.charts-submanifold } C)$  for  $x$ 
  proof –
  have  $X p ?f' = \text{vec-field-apply-fun-in-at } X f U p$  if  $p: p \in C$  for  $p$ 
  using  $U.\text{vec-field-apply-fun-in-open}[OF U(2) C(2) - - 1 f - X, \text{symmetric}]$ 
  using  $C p f f'(2) C(4)$  by (auto simp: subset-iff)
  thus thesis using that by simp
  qed
  qed
  thus  $\text{diff} \infty (U.\text{sub.charts-submanifold } C) \text{charts-eucl} (\text{vec-field-apply-fun-in-at}$ 
 $X f U)$ 

```

```

    unfolding diff-fun-def .
  qed (simp-all add: C(1,2))
qed
qed

moreover have ?LHS if diff-apply-fun: ?RHS2
proof (intro smooth-vector-field-if-local ballI)
  fix p assume p: p ∈ carrier
  obtain c where c: c ∈ atlas p ∈ domain c using p atlasE by blast
  then interpret p: c-manifold-point charts ∞ c p using c-manifold-point by
blast

  have vf-smooth-local-iff: smooth-vector-field-local charts c X
  if (∀ i ∈ Basis. diff-fun-on (domain c) (p.vector-field-component X i))
  using p.vector-field-smooth-local-iff rough-vector-field-subset[OF X]
  by (meson assms in-carrier-atlasI p.ψ rough-vector-field-def that)
  have smooth-vector-field-local charts c X
  proof (rule vf-smooth-local-iff, intro ballI)
    fix i::'b assume i: i ∈ Basis
    — Add simp rules to make diff-fun-on equivalent to diff-fun on domain c.
    note local-simps[simp] = diff-fun-on-open domain-atlas-subset-carrier

    define apply-X-in-c-at (↪ X_c'' -> [80,80])
    where [simp]: apply-X-in-c-at ≡ λf q. vec-field-apply-fun-in-at X f (domain
c) q

    have (X|(domain c)) q (p.local-coord-at p i) = p.vector-field-component X i q
    if q ∈ domain c for q
    proof —
      interpret q: c-manifold-point charts ∞ c q using that c by unfold-locales
simp-all
      have Xq: X q ∈ q.T_p M using rough-vector-fieldE[OF X] that c(1) by blast
      show ?thesis
        using vec-field-apply-fun-alt[OF open-domain that] apply (simp add:
p.local-coord-diff-fun X)
        using restrict0-apply-in[OF Xq]
        using p.vector-apply-coord-at p.p that q.component-function-apply-in-T_p M
        by (smt (verit, best) Xq assms c-manifold-local.vec-field-apply-fun-eq-component
p.c-manifold-local-axioms p.local-coord-diff-fun)
    qed
    moreover have diff-fun-on (domain c) (λq. X_c'' (p.local-coord-at p i) q)
    using diff-apply-fun by (simp add: p.local-coord-diff-fun)

    ultimately show diff-fun-on (domain c) (p.vector-field-component X i)
    by (simp add: diff-fun-on-def)
  qed
with c show ∃ c ∈ atlas. p ∈ domain c ∧ smooth-vector-field-local charts c X by
blast

```


qed (*simp add: assms rough-vector-fieldE(2)*)

ultimately show $?LHS \longleftrightarrow ?RHS1$ $?LHS \longleftrightarrow ?RHS2$ **by auto**
qed

lemma *vector-field-smooth-iff'*:

fixes *C-inf*

defines $\bigwedge U. C\text{-inf } U \equiv c\text{-manifold.diff-fun-space } (charts\text{-submanifold } U) \infty$

assumes *X: rough-vector-field X*

shows *smooth-vector-field X* $\longleftrightarrow (\forall f \in \text{diff-fun-space}. (X'' f) \in \text{diff-fun-space})$

and *smooth-vector-field X* $\longleftrightarrow (\forall U f. \text{open } U \wedge U \subseteq \text{carrier} \wedge f \in C\text{-inf } U$

\longrightarrow

$\text{diff-fun-on } U (X \upharpoonright U'' f))$

proof –

show *smooth-vector-field X* = $(\forall U f. \text{open } U \wedge U \subseteq \text{carrier} \wedge f \in C\text{-inf } U \longrightarrow \text{diff-fun-on } U (\lambda p. X \upharpoonright U p f))$

apply (*simp add: diff-fun-on-open assms(1)*)

using *vector-field-smooth-iff(2)[OF assms(2–)]*

by (*smt (verit, ccfv-SIG) Un-Int-eq(4) diff-fun.diff-fun-cong open-imp-submanifold restrict0-apply-in submanifold.carrier-submanifold sup.order-iff*)

qed (*simp add: vector-field-smooth-iff(1)[OF assms(2–)]*)

lemma *smooth-vf-diff-fun-space*:

assumes *X: smooth-vector-field X*

and *f: f ∈ diff-fun-space*

shows $Xf \in \text{diff-fun-space}$

using *vector-field-smooth-iff(1) smooth-vector-field-def X f rough-vector-fieldE*

by (*auto simp: diff-fun-space-def extensional0-def*)

end

5.4 Smooth vector fields are derivations

context *c-manifold begin*

— Generalising *is-derivation* (which might have been called *is-derivation-at*) over the carrier set. Relative to that definition, we also add a condition on the codomain.

definition *is-derivation-on* :: $((\text{'a} \Rightarrow \text{real}) \Rightarrow (\text{'a} \Rightarrow \text{real})) \Rightarrow \text{bool}$ **where**

is-derivation-on D $\equiv \text{real-linear-on } \text{diff-fun-space } \text{diff-fun-space } D \wedge$

$(\forall f \in \text{diff-fun-space}. \forall g \in \text{diff-fun-space}. D (f * g) = f * (D g) +$

$g * (D f)) \wedge$

$D \text{ ' diff-fun-space } \subseteq \text{diff-fun-space}$

lemma *vec-field-linear-on*:

assumes *X: rough-vector-field X*

and *b: b1 ∈ diff-fun-space b2 ∈ diff-fun-space*

shows $X'' (b1+b2) = (Xb1 + Xb2) X'' (r *_R b1) = (r *_R (Xb1))$
using *tangent-space-linear-on*[*OF rough-vector-fieldE(1)*][*OF X*]
by (*auto simp: linear-on-def module-hom-on-def module-hom-on-axioms-def assms(2-)*)

lemma *linear-on-vec-field*:
assumes *rough-vector-field X*
shows *real-linear-on diff-fun-space diff-fun-space* ($() X$)
using *vec-field-linear-on*[*OF assms*] **by** (*unfold-locales, auto*)

lemma *product-rule-vf*:
assumes *X: rough-vector-field X*
and $f \in \text{diff-fun-space } g \in \text{diff-fun-space}$
shows $X'' (f*g) = f * (X'' g) + g * (X'' f)$
using *tangent-space-derivation rough-vector-fieldE(1) assms* **by** *auto*

end

context *smooth-manifold* **begin**

lemma *vector-field-is-derivation*:
assumes *X: smooth-vector-field X*
shows *is-derivation-on* ($\lambda f. Xf$)
using *linear-on-vec-field product-rule-vf vector-field-smooth-iff(1)*
using *smooth-vector-field-def assms unfolding is-derivation-on-def* **by** *auto*

5.5 Derivations are smooth vector fields

lemma *extensional-derivation-is-smooth-vector-field*:
fixes $D :: ('a \Rightarrow \text{real}) \Rightarrow ('a \Rightarrow \text{real})$ **and** $X :: 'a \Rightarrow ('a \Rightarrow \text{real}) \Rightarrow \text{real}$
defines [*simp*]: $X \equiv \lambda p. \lambda f. D f p$
assumes *der-D: is-derivation-on D*
and *ext-X: extensional0 carrier X*
and *ext-D: extensional0 diff-fun-space D*
shows *smooth-vector-field X*

proof –

have *rough-vf-X: rough-vector-field X*
proof (*unfold rough-vector-field-def tangent-space-def is-derivation-def, safe*)
fix p **assume** $p: p \in \text{carrier}$
show *extensional0 diff-fun-space* ($\lambda f. X p f$)
by (*simp, metis (mono-tags, lifting) ext-D extensional0-def zero-fun-apply*)
show *real-linear-on diff-fun-space UNIV* ($\lambda f. X p f$)
using *der-D[unfolded is-derivation-on-def]*
unfolding *linear-on-def module-hom-on-def module-hom-on-axioms-def*
using *manifold-eucl.diff-fun-space.m2.module-on-axioms* **by** *auto*
fix $f g$ **assume** $f \in \text{diff-fun-space } g \in \text{diff-fun-space}$
thus $X p (f * g) = f p * X p g + g p * X p f$
using *der-D[unfolded is-derivation-on-def]* **by** *simp*
qed (*rule ext-X*)

show *smooth-vector-field* X
unfolding *vector-field-smooth-iff*(1)[*OF rough-vf-X*]
using *der-D*[*unfolded is-derivation-on-def*] *diff-fun-spaceD* X -*def* **by** *blast*
qed

lemma *extensional-derivation-is-smooth-vector-field'*:
fixes $D :: ('a \Rightarrow \text{real}) \Rightarrow ('a \Rightarrow \text{real})$
assumes *der-D*: *is-derivation-on* D
and *ext-X*: *extensional0 carrier* $(\lambda p f. D f p)$
and *ext-D*: *extensional0 diff-fun-space* D
obtains X **where** *smooth-vector-field* X **and** $\forall f \in \text{diff-fun-space}. D f = X f$
using *extensional-derivation-is-smooth-vector-field*[*OF assms*] **by** *auto*

theorem *smooth-vector-field-iff-derivation*:
fixes *extensional-derivation* **defines** $\bigwedge D. \text{extensional-derivation } D \equiv$
is-derivation-on $D \wedge \text{extensional0 carrier } (\lambda p f. D f p) \wedge \text{extensional0 diff-fun-space}$
 D
shows *smooth-vector-field* $X \implies \text{extensional-derivation } (\lambda f. X'' f)$
and *extensional-derivation* $D \implies \text{smooth-vector-field } (\lambda p f. D f p)$
unfolding *assms*(1) **using** *vector-field-is-derivation extensional-derivation-is-smooth-vector-field*
by (*simp-all add: ext0-vec-field-apply-fun smooth-vector-fieldE*(2,3))

end

end

6 The Lie bracket of smooth vector fields

theory *Manifold-Lie-Bracket*

imports

Smooth-Vector-Fields

Algebra-On

begin

definition *lie-bracket-of-smooth-vector-fields* :: *'a vector-field* \Rightarrow *'a vector-field* \Rightarrow
'a vector-field

where *lie-bracket-of-smooth-vector-fields* $X Y \equiv \lambda p :: 'a. \lambda f :: 'a \Rightarrow \text{real}. X p (Y'' f) - Y p (X'' f)$

notation *lie-bracket-of-smooth-vector-fields* ($\langle [-; -] \rangle$ [65,65])

lemma *lie-bracket-def*: $[X; Y] p f = X p (Y f) - Y p (X f)$

unfolding *lie-bracket-of-smooth-vector-fields-def* **by** *simp*

context *c-manifold* **begin**

6.1 General lemmas

lemma *is-derivation-uminus*: *is-derivation* $(-x) p$ **if** x : *is-derivation* $x p$

using *is-derivation-scaleR*[*OF* *x*, *of* -1] **by** *simp*

lemma *is-derivation-minus*: *is-derivation* $(x - y)$ *p*
if *x*: *is-derivation* *x* *p* **and** *y*: *is-derivation* *y* *p*
using *is-derivation-add*[*OF* *x* *is-derivation-uminus*[*OF* *y*]] **by** *simp*

lemma *diff-fun-space-minus*: $f - g \in \text{diff-fun-space}$
if $f \in \text{diff-fun-space}$ $g \in \text{diff-fun-space}$
by (*simp* *add*: *diff-fun-space.m1.subspace-UNIV* *diff-fun-space.m1.subspace-diff*
that(1) *that*(2))

lemma *rough-vector-field-add*:
assumes *rough-vector-field* *X* *rough-vector-field* *Y*
shows *rough-vector-field* $(X + Y)$
using *assms* *rough-vector-field-def* *tangent-space.mem-add* **by** *force*

abbreviation (*input*) *scaleR-vf* \equiv *scaleR* :: *real* \Rightarrow '*a* *vector-field* \Rightarrow '*a* *vector-field*

lemma *scaleR-vf*: *scaleR-vf* $= (\lambda r$ *X* *p* *f*. $r * X$ *p* *f*) **by** *fastforce*

lemma *rough-vector-field-scaleR*:
assumes *rough-vector-field* *X*
shows *rough-vector-field* $(\text{scaleR-vf } a$ *X*)
using *assms* *tangent-space.mem-scale* **by** (*simp* *add*: *rough-vector-field-def*)

6.2 Properties of the Lie bracket on \mathfrak{X}

lemma *lie-bracket-antisym*: $[X; Y] = -[Y; X]$
unfolding *lie-bracket-def* **by** *fastforce*

lemma *ext0-lie-bracket*:

shows *extensional0* *carrier* *X* \implies *extensional0* *carrier* *Y* \implies *extensional0* *carrier*
 $[X; Y]$

and *rough-vector-field* *X* \implies *rough-vector-field* *Y* \implies *extensional0* *diff-fun-space*
 $(\text{vec-field-apply-fun } [X; Y])$

proof –

show *extensional0* *carrier* *X* \implies *extensional0* *carrier* *Y* \implies *extensional0* *carrier*
 $[X; Y]$

unfolding *lie-bracket-def* *extensional0-def* **by** *auto*

assume *asm*: *rough-vector-field* *X* *rough-vector-field* *Y*

then show *extensional0* *diff-fun-space* $(\text{vec-field-apply-fun } [X; Y])$

proof – — This proof was fiddly to get into a form where methods would not time out.

note *vf-0* $=$ *linear-on.linear-0*[*OF* *linear-on-vec-field*]

have $\forall p$. $(X'' 0)$ *p* $- (Y'' 0)$ *p* $= 0$

using *vf-0*[*OF* *asm*(1)] *vf-0*[*OF* *asm*(2)] **by** (*metis* *diff-0-right* *zero-fun-apply*)

then have 0 : $(\lambda p$. $(X'' 0)$ *p* $- (Y'' 0)$ *p*) $= 0$ **by** *auto*

```

    thus ?thesis
      using ext0-vec-field-apply-fun asm unfolding lie-bracket-def extensional0-def
    by presburger
  qed
end

```

context *smooth-manifold* **begin**

A nice computational proof that I try to keep close-ish to Lee's original pen-and-paper [?, p. 186].

lemma *product-rule-lie-bracket*:

```

  assumes X: smooth-vector-field X
    and Y: smooth-vector-field Y
    and diff-funs: f ∈ diff-fun-space g ∈ diff-fun-space
  shows [X; Y]" (f * g) = f * [X; Y]" g + g * [X; Y]" f
  proof -
    have rough-vf[simp]: rough-vector-field X rough-vector-field Y
      using smooth-vector-field-def assms by auto
    interpret linear-X: linear-on diff-fun-space diff-fun-space scaleR scaleR vec-field-apply-fun X
      using linear-on-vec-field[OF rough-vf(1)] .
    interpret linear-Y: linear-on diff-fun-space diff-fun-space scaleR scaleR vec-field-apply-fun Y
      using linear-on-vec-field[OF rough-vf(2)] .
  
```

note [simp] = *diff-funs diff-fun-space.m1.mem-add diff-fun-space-times*

have *local-simps* [simp]:

```

  ∧ f. f ∈ diff-fun-space ⇒ X" f ∈ diff-fun-space
  ∧ f. f ∈ diff-fun-space ⇒ Y" f ∈ diff-fun-space
  using X Y by (simp-all add: vector-field-smooth-iff(1))

```

have $([X; Y]" (f * g)) = X" (Y(f * g)) - Y" (X(f * g))$

unfolding *lie-bracket-def* **by** (*simp add: fun-diff-def*)

also have $\dots = X" (f * Yg + g * Yf) - Y" (f * Xg + g * Xf)$

using *rough-vf diff-funs product-rule-vf* **by** *presburger*

also have $\dots = X" (f * Yg) + X" (g * Yf) - Y" (f * Xg) - Y" (g * Xf)$

— Extra step to invoke linearity of both X and Y.

using *linear-X.add linear-Y.add* **by** *simp*

also have $\dots = (f * X(Yg)) + (g * X(Yf)) - (f * Y(Xg)) - (g * Y(Xf))$

— No separate step for term cancellation.

using *product-rule-vf* **by** *auto*

finally show *?thesis*

```

  by (simp add: lie-bracket-def fun-diff-def add-diff-eq diff-add-eq plus-fun-def
    vector-space-over-itself.scale-right-diff-distrib)

```

qed

lemma *lie-bracket-is-derivation-on*:
assumes X : *smooth-vector-field* X
and Y : *smooth-vector-field* Y
shows *is-derivation-on* $(\lambda f. [X; Y]'' f)$
proof (*unfold is-derivation-on-def, safe*)
have 0 : $(\lambda p. Y p (\lambda p. X p f)) \in \text{diff-fun-space } (\lambda p. X p (\lambda p. Y p f)) \in \text{diff-fun-space}$
if f : $f \in \text{diff-fun-space}$ **for** f
using *smooth-vf-diff-fun-space*[OF Y *smooth-vf-diff-fun-space*[OF X f]]
using *smooth-vf-diff-fun-space*[OF X *smooth-vf-diff-fun-space*[OF Y f]] **by**
simp-all
show 1 : $[X; Y]'' f \in \text{diff-fun-space}$ **if** f : $f \in \text{diff-fun-space}$ **for** f
using 0 [OF f] *diff-fun-space-minus* **by** (*simp add: fun-diff-def lie-bracket-def*)
thus 2 : $[X; Y]'' (f * g) = f * ([X; Y]'' g) + g * ([X; Y]'' f)$
if f : $f \in \text{diff-fun-space}$ **and** g : $g \in \text{diff-fun-space}$ **for** f g
using *product-rule-lie-bracket*[OF X Y f g] **by** *simp*
show 3 : *linear-on diff-fun-space diff-fun-space scaleR scaleR* $(\lambda f. [X; Y]'' f)$
proof –
have $lin-X$: *real-linear-on diff-fun-space diff-fun-space* $(\lambda f. X f)$
using *linear-on-vec-field* X [*unfolded smooth-vector-field-def*] **by** *simp*
have $lin-Y$: *real-linear-on diff-fun-space diff-fun-space* $(\lambda f. Y f)$
using *linear-on-vec-field* Y [*unfolded smooth-vector-field-def*] **by** *simp*
have $lin-XY$: *real-linear-on diff-fun-space diff-fun-space* $((\text{vec-field-apply-fun } X)$
 $\circ (\text{vec-field-apply-fun } Y))$
using *smooth-vf-diff-fun-space*[OF Y] **by** (*auto intro: linear-on-compose*[OF
 $lin-Y$ $lin-X$])
have $lin-YX$: *real-linear-on diff-fun-space diff-fun-space* $((\text{vec-field-apply-fun } Y)$
 $\circ (\text{vec-field-apply-fun } X))$
using *smooth-vf-diff-fun-space*[OF X] **by** (*auto intro: linear-on-compose*[OF
 $lin-X$ $lin-Y$])
have *real-linear-on diff-fun-space diff-fun-space* $(\lambda x. (X'' (Y'' x)) - (Y'' (X'' x)))$
apply (*intro vector-space-pair-on.linear-compose-sub*[OF - - - $lin-XY$ $lin-YX$,
simplified])
using *linear-on.vector-space-pair-on*[OF $lin-X$] 0 **by** *auto*
then show *?thesis* **by** (*simp add: lie-bracket-def fun-diff-def*)
qed
qed

This is Lee's [?, Lemma 8.25].

lemma *lie-bracket-closed*:
assumes X : *smooth-vector-field* X
and Y : *smooth-vector-field* Y
shows *smooth-vector-field* $[X; Y]$
using *extensional-derivation-is-smooth-vector-field lie-bracket-is-derivation-on*
ext0-lie-bracket assms smooth-vector-field-def rough-vector-fieldE(2) **by** *auto*

lemma

```

assumes  $X$ : smooth-vector-field  $X$ 
and  $Y$ : smooth-vector-field  $Y$ 
and  $Z$ : smooth-vector-field  $Z$ 
shows lie-bracket-add-left:  $[X+Y;Z] = [X;Z] + [Y;Z]$ 
and lie-bracket-add-right:  $[X;Y+Z] = ([X;Y] + [X;Z])$ 
proof -
have distrib-left:  $[X+Y;Z] = ([X;Z] + [Y;Z])$ 
if  $X$ : smooth-vector-field  $X$ 
and  $Y$ : smooth-vector-field  $Y$ 
and  $Z$ : smooth-vector-field  $Z$ 
for  $X Y Z$ 
proof (standard+)
fix  $p f$ 
show  $[X+Y;Z] p f = ([X;Z] + [Y;Z]) p f$ 
proof (cases  $p \in \text{carrier} \wedge f \in \text{diff-fun-space}$ )

```

We deal with the cases outside our interest off the bat. This is just taking care of *extensional0* in both (point and function) arguments of the vector field.

```

case False
then show ?thesis
apply (cases  $p \in \text{carrier}$ , simp-all)
subgoal
using False  $X Y Z$  smooth-vector-field-def rough-vector-field-add[of  $X Y$ ]
using extensional0-outside[OF - ext0-lie-bracket(2)]
extensional0-add[OF smooth-vector-fieldE(2)[OF  $X$ ] smooth-vector-fieldE(2)[OF
 $Y$ ]]
by (smt (verit, ccfv-SIG) zero-fun-apply)
using  $X Y Z$  smooth-vector-fieldE(2) extensional0-outside[OF - ext0-lie-bracket(1)]
by force
next

```

The rest of this proof is just linearity of the tangent vector $Z p$.

```

case True hence  $p$ :  $p \in \text{carrier}$  and  $f$ :  $f \in \text{diff-fun-space}$  by simp+
interpret linZ: linear-on diff-fun-space UNIV scaleR scaleR  $Z p$ 
using tangent-spaceD(1)[OF smooth-vector-fieldE(1)[OF  $Z$ ]] by blast
show ?thesis
using linZ.add  $X Y$  smooth-vf-diff-fun-space  $f$  by (auto simp: lie-bracket-def
plus-fun-def)
qed
qed
thus  $[X+Y;Z] = [X;Z] + [Y;Z]$  using assms by blast
show  $[X;Y+Z] = ([X;Y] + [X;Z])$ 
using distrib-left[OF  $Y Z X$ ] lie-bracket-antisym by (metis minus-add-distrib)
qed

```

lemma

```

assumes  $X$ : smooth-vector-field  $X$ 

```

and Y : *smooth-vector-field* Y
shows *lie-bracket-scale-left*: $[scaleR\text{-}vf\ a\ X; Y] = scaleR\text{-}vf\ a\ [X; Y]$
and *lie-bracket-scale-right*: $[X; scaleR\text{-}vf\ a\ Y] = scaleR\text{-}vf\ a\ [X; Y]$
proof –

We proceed as above, dealing with extensionality before using an existing linearity result.

have *scaleR-vf-left*: $[scaleR\text{-}vf\ a\ X; Y] = scaleR\text{-}vf\ a\ [X; Y]$
if X : *smooth-vector-field* X
and Y : *smooth-vector-field* Y
for $X\ Y\ a$
proof (*standard+*)
fix $p\ f$
show $[scaleR\text{-}vf\ a\ X; Y]\ p\ f = scaleR\text{-}vf\ a\ [X; Y]\ p\ f$
proof (*cases* $p \in carrier \wedge f \in diff\text{-}fun\text{-}space$)
case *False*
then show *?thesis*
apply (*cases* $p \in carrier$)
subgoal
using *False* $X\ Y$ *smooth-vector-field-def* *rough-vector-field-scaleR*
using *extensional0-outside*[*OF* - *ext0-lie-bracket*(2)] *extensional0-scaleR*
by (*smt* (*verit*, *del-insts*) *scaleR-cancel-right* *scaleR-fun-beta* *scaleR-zero-left*)
using *smooth-vector-fieldE*(2) $X\ Y$ *extensional0-outside*[*OF* - *ext0-lie-bracket*(1)]
by *simp*
next
case *True* **hence** f : $f \in diff\text{-}fun\text{-}space$ **by** *simp+*
interpret $linY$: *linear-on* *diff-fun-space* *UNIV* *scaleR* *scaleR* $Y\ p$
using *tangent-spaceD*(1)[*OF* *smooth-vector-fieldE*(1)[*OF* Y]] **by** *blast*
show *?thesis*
using $linY$.*scale*[*OF* *smooth-vf-diff-fun-space*, *OF* $X\ f$]
by (*auto* *simp*: *lie-bracket-def* *scaleR-fun-def* *right-diff-distrib*)
qed
qed
thus $[scaleR\text{-}vf\ a\ X; Y] = scaleR\text{-}vf\ a\ [X; Y]$ **by** (*simp* *add*: $X\ Y$)
show $[X; scaleR\text{-}vf\ a\ Y] = scaleR\text{-}vf\ a\ [X; Y]$
apply (*simp* *only*: *lie-bracket-antisym*[*of* X *scaleR-vf* $a\ Y$] *lie-bracket-antisym*[*of* $X\ Y$])
using *scaleR-vf-left*[*OF* $Y\ X$] **by** *fastforce*
qed

lemmas *lie-bracket-bilinear-simps* [*simp*] = *lie-bracket-scale-left*
lie-bracket-scale-right
lie-bracket-add-left
lie-bracket-add-right

lemma (*in* *module-hom-on*) *diff*:
 $b1 \in S1 \implies b2 \in S1 \implies f\ (b1 - b2) = f\ b1 - f\ b2$
by (*metis* *add* *diff-eq-eq* *m1.subspace-UNIV* *m1.subspace-diff* *set-eq-subset*)


```

lemma lie-bracket-jacobi:  $[X; [Y;Z]] + [Y;[Z;X]] + [Z;[X;Y]] = 0$ 
  if  $X$ : smooth-vector-field  $X$ 
    and  $Y$ : smooth-vector-field  $Y$ 
    and  $Z$ : smooth-vector-field  $Z$ 
proof –
  have rough-vf: rough-vector-field  $X$  rough-vector-field  $Y$  rough-vector-field  $Z$ 
    using smooth-vector-field-def that by auto
  interpret linear-X: linear-on diff-fun-space diff-fun-space scaleR scaleR vec-field-apply-fun
 $X$ 
    using linear-on-vec-field[OF rough-vf(1)] .
  interpret linear-Y: linear-on diff-fun-space diff-fun-space scaleR scaleR vec-field-apply-fun
 $Y$ 
    using linear-on-vec-field[OF rough-vf(2)] .
  interpret linear-Z: linear-on diff-fun-space diff-fun-space scaleR scaleR vec-field-apply-fun
 $Z$ 
    using linear-on-vec-field[OF rough-vf(3)] .

  have local-simps:
     $\wedge f. f \in \text{diff-fun-space} \implies X'' f \in \text{diff-fun-space}$ 
     $\wedge f. f \in \text{diff-fun-space} \implies Y'' f \in \text{diff-fun-space}$ 
     $\wedge f. f \in \text{diff-fun-space} \implies Z'' f \in \text{diff-fun-space}$ 
    using  $X Y Z$  by (simp-all add: vector-field-smooth-iff rough-vf)

  {
    fix  $f$  assume  $f: f \in \text{diff-fun-space}$ 
    have  $[X; [Y;Z]]'' f + [Y;[Z;X]]'' f + [Z;[X;Y]]'' f =$ 
       $X'' ([Y;Z]f) - [Y;Z]'' (Xf) + Y'' ([Z;X]f) - [Z;X]'' (Yf) + Z'' ([X;Y]f) -$ 
 $[X;Y]'' (Zf)$ 
    unfolding lie-bracket-def by auto
    also have  $\dots = X(Y(Zf)) - X(Z(Yf))$ 
       $- Y(Z(Xf)) + Z(Y(Xf))$ 
       $+ Y(Z(Xf)) - Y(X(Zf))$ 
       $- Z(X(Yf)) + X(Z(Yf))$ 
       $+ Z(X(Yf)) - Z(Y(Xf))$ 
       $- X(Y(Zf)) + Y(X(Zf))$ 
    using linear-X.diff linear-Y.diff linear-Z.diff by (auto simp: f fun-diff-def
lie-bracket-def local-simps)
    finally have  $[X; [Y;Z]]'' f + [Y;[Z;X]]'' f + [Z;[X;Y]]'' f = 0$  by simp
  } moreover {
    fix  $f$  assume  $f: f \notin \text{diff-fun-space}$ 
    have  $[X; [Y;Z]]'' f = 0$   $[Y; [Z;X]]'' f = 0$   $[Z;[X;Y]]'' f = 0$ 
    using ext0-lie-bracket(2)[OF smooth-vector-fieldE(3)] smooth-vector-fieldE(3)
 $X Y Z$ 
    using lie-bracket-closed(1)[OF Y Z] lie-bracket-closed(1)[OF Z X] lie-bracket-closed(1)[OF
 $X Y$ ]
    by (simp-all add: extensional0-def f smooth-vector-field-alt)
    hence  $[X; [Y;Z]]'' f + [Y;[Z;X]]'' f + [Z;[X;Y]]'' f = 0$  by simp

```

} ultimately have $\bigwedge p f. [X; [Y; Z]] p f + [Y; [Z; X]] p f + [Z; [X; Y]] p f = 0$
by (*smt (verit, best) plus-fun-apply zero-fun-apply*)
thus ?thesis by (*intro HOL.ext*) *fastforce*
qed

definition $SVF \equiv \{X. \text{smooth-vector-field } X\}$

lemma *lie-algebra-of-smooth-vector-fields: lie-algebra SVF scaleR-vf lie-bracket-of-smooth-vector-fields*
proof –

note *svf-if-derivI = extensional-derivation-is-smooth-vector-field[unfolded is-derivation-on-def]*

have *svf-0: smooth-vector-field 0*
apply (*intro svf-if-derivI, safe, unfold-locales*)
apply *auto[3]*
using *diff-fun-space.m1.mem-zero uminus-apply* **apply** *fastforce*
using *extensional0-def zero-fun-def* **by** *auto*

have *local-simps: ($\lambda r f p. r * f (p::'a) = \text{scaleR}$)*
by *fastforce*

have *svf-scaleR: smooth-vector-field (a *_R X)*
if *X: smooth-vector-field X for a X*

proof (*intro svf-if-derivI, intro conjI ballI*)

show *extensional0 carrier (a *_R X) by (simp add: smooth-vector-fieldE(2) that)*

have *derX: linear-on diff-fun-space diff-fun-space scaleR scaleR ($\lambda f p. X p f$)*
 $(\bigwedge g. f \in \text{diff-fun-space} \implies g \in \text{diff-fun-space} \implies X'' (f * g) = f * (X'' g) + g * (X'' f))$

$(\lambda f p. X p f) \text{ ' diff-fun-space} \subseteq \text{diff-fun-space}$

using *X vector-field-is-derivation unfolding is-derivation-on-def* **by** *auto*

show *real-linear-on diff-fun-space diff-fun-space ($\lambda f p. (a *_{\mathbb{R}} X) p f$)*

using *linear-on-compose[OF derX(1) diff-fun-space.m1.linear-scale-self derX(3)]*

by (*simp add: o-def, metis local-simps*)

show $(\lambda p. (a *_{\mathbb{R}} X) p (f * g)) = f * (\lambda p. (a *_{\mathbb{R}} X) p g) + g * (\lambda p. (a *_{\mathbb{R}} X) p f)$

if *f ∈ diff-fun-space g ∈ diff-fun-space for f g*

proof –

have $(\lambda p. a * X p (f * g)) = (\lambda x. a) * (\lambda p. X p (f * g))$ **by** *auto*

then show $(\lambda p. (a *_{\mathbb{R}} X) p (f * g)) = f * (\lambda p. (a *_{\mathbb{R}} X) p g) + g * (\lambda p. (a *_{\mathbb{R}} X) p f)$

using *derX(2)[OF that]* **by** (*auto simp: distrib-left*)

qed

show $(\lambda f p. (a *_{\mathbb{R}} X) p f) \text{ ' diff-fun-space} \subseteq \text{diff-fun-space}$

using *derX(3) diff-fun-space.m1.mem-scale* **by** (*auto, metis image-subset-iff local-simps*)

show *extensional0 diff-fun-space ($\lambda f p. (a *_{\mathbb{R}} X) p f$)*

```

    using X[unfolded smooth-vector-field-def] ext0-vec-field-apply-fun
    by (meson extensional0-scaleR rough-vector-field-scaleR)
qed

have svf-add: smooth-vector-field (X + Y)
  if X: smooth-vector-field X and Y: smooth-vector-field Y
  for X Y
proof (intro svf-if-derivI, intro conjI ballI)
  have derX: real-linear-on diff-fun-space diff-fun-space (λf p. X p f)
    (λf g. f ∈ diff-fun-space ⇒ g ∈ diff-fun-space ⇒ X'' (f * g) = f * (X'' g) +
  g * (X'' f))
    (λf p. X p f) ' diff-fun-space ⊆ diff-fun-space
  and derY: real-linear-on diff-fun-space diff-fun-space (λf p. Y p f)
    (λf g. f ∈ diff-fun-space ⇒ g ∈ diff-fun-space ⇒ Y'' (f * g) = f * (Y'' g) +
  g * (Y'' f))
    (λf p. Y p f) ' diff-fun-space ⊆ diff-fun-space
  using X Y vector-field-is-derivation unfolding is-derivation-on-def by auto

interpret D: vector-space-pair-on diff-fun-space diff-fun-space scaleR scaleR by
unfold-locales

show real-linear-on diff-fun-space diff-fun-space (λf p. (X + Y) p f)
  apply (simp, intro D.linear-compose-add[unfolded plus-fun-def])
  using derX(1,3) derY(1,3) by auto
show (λp. (X + Y) p (f * g)) = f * (λp. (X + Y) p g) + g * (λp. (X + Y)
p f)
  if f ∈ diff-fun-space g ∈ diff-fun-space for f g
  using derX(2)[OF that] derY(2)[OF that]
  by (simp add: plus-fun-def distrib-left, metis (no-types) add commute add.left-commute)
show (λf p. (X + Y) p f) ' diff-fun-space ⊆ diff-fun-space
  using derX(3) derY(3) diff-fun-space.m1.mem-add by (auto simp: plus-fun-def)
show extensional0 diff-fun-space (λf p. (X + Y) p f)
  using X Y ext0-vec-field-apply-fun rough-vector-field-add smooth-vector-field-def
by blast
show extensional0 carrier (X + Y) by (simp add: X Y smooth-vector-fieldE(2))
qed

interpret vector-space-svf: vector-space-on SVF scaleR-vf
  using svf-0 svf-scaleR svf-add SVF-def
  by (unfold-locales, auto simp: scaleR-right-distrib scaleR-left-distrib)

have lie-bracket-antisym': [X;X] = 0
  if X: smooth-vector-field X extensional0 carrier X for X
  using lie-bracket-antisym by (metis one-neq-neg-one scaleR-cancel-right scaleR-minus1-left
scaleR-one)

show ?thesis
  apply (intro vector-space-svf.lie-algebraI, unfold SVF-def)
  using lie-bracket-closed lie-bracket-antisym' lie-bracket-jacobi

```

by (*simp-all add: smooth-vector-fieldE(2)*)
qed

end

end

theory *Lie-Group*

imports

HOL-Analysis.Analysis

HOL-Eisbach.Eisbach

More-Manifolds

begin

7 Definition of Lie Groups (as Locales)

Some abbreviations for easier reading first. A binary operation is colloquially said continuous/smooth/differentiable on a manifold M if it is so on the product manifold M^2 . We fix the types of the binary operations in two of the definitions below, as the target space is made explicit only in the third (the one using *diff* ∞).

abbreviation (*input*) *continuous-on-product-manifold charts* (*binop*:: $'a \Rightarrow 'a \Rightarrow 'a$:: $\{second-countable-topology, t2\}$)
 \equiv

continuous-on (*c-manifold-prod.carrier charts charts*) ($\lambda(a,b)$. *binop* *a b*)

abbreviation (*input*) *smooth-on-product-manifold charts* (*binop*:: $'a \Rightarrow 'a \Rightarrow 'a$:: $\{second-countable-topology, real-n\}$)
 \equiv

smooth-on (*c-manifold-prod.carrier charts charts*) ($\lambda(a,b)$. *binop* *a b*)

abbreviation (*input*) *diff-on-product-manifold charts* *binop* \equiv

diff ∞ (*c-manifold-prod.prod-charts charts charts*) *charts* ($\lambda(a,b)$. *binop* *a b*)

7.1 Topological groups

A group with a topology, such that the group operations are continuous.

locale *topological-group* =

manifold charts + *group-on-with carrier tms tms-one dvsn invs*

for *charts*::($'a$:: $\{t2-space, second-countable-topology\}$, $'e$::*euclidean-space*) *chart set*

and *tms tms-one dvsn invs* +

assumes *cts-mult: continuous-on-product-manifold charts tms*

and *cts-inv: continuous-on carrier invs*

7.2 Lie groups

A Lie group is a group on a set, but instead of a carrier set, we specify a set of charts, which imply the carrier set as a (smooth) manifold M . Internally, we consider the product manifold, to define smoothness of multiplication

$M \times M \rightarrow M$. It may be overkill to keep inverse and division separate, considering *group-on-with* includes an axiom to relate the two, but this is how it's done in other Isabelle theories, so I'll keep it. It gives some extra flexibility, and an intro lemma using the more traditional group parameters (an operation, and an identity) and axioms is already provided in $\llbracket \forall a \in ?G. \forall b \in ?G. ?mult\ a\ b \in ?G; \forall a \in ?G. \forall b \in ?G. \forall c \in ?G. ?mult\ (?mult\ a\ b)\ c = ?mult\ a\ (?mult\ b\ c); ?e \in ?G \wedge (\forall a \in ?G. ?mult\ ?e\ a = a \wedge ?mult\ a\ ?e = a); \forall x \in ?G. \exists y. y \in ?G \wedge ?mult\ x\ y = ?e \wedge ?mult\ y\ x = ?e \rrbracket \implies \text{group-on-with } ?G\ ?mult\ ?e\ (\lambda x\ z. ?mult\ x\ (THE\ y. y \in ?G \wedge ?mult\ z\ y = ?e \wedge ?mult\ y\ z = ?e))\ (\lambda x. THE\ y. y \in ?G \wedge ?mult\ x\ y = ?e \wedge ?mult\ y\ x = ?e)$.

locale *lie-group* =
c-manifold charts ∞ + *group-on-with carrier tms tms-one dvn invs*
for *charts::('a::(t2-space,second-countable-topology), 'e::euclidean-space) chart set*
and *tms tms-one dvn invs* +
assumes *smooth-mult: diff-on-product-manifold charts tms*
and *smooth-inv: diff ∞ charts charts invs*

We can make a shortened locale for Lie groups where the inversion and division are implied. This does *not* say anything about the implementation of inversion or division outside the carrier set. See also *grp-on*.

locale *lie-grp* =
c-manifold charts ∞ + *grp-on carrier tms one*
for *charts::('a::(t2-space,second-countable-topology), 'e::euclidean-space) chart set*
and *tms one* +
— multiplication and inversion are smooth
assumes *smooth-mult: diff-on-product-manifold charts tms*
and *smooth-inv: diff ∞ charts charts invs*
begin

lemma *is-lie-group: lie-group charts tms one mns invs*
unfolding *lie-group-def lie-group-axioms-def*
by (*auto simp: c-manifold-axioms smooth-mult is-group-on-with smooth-inv*)

sublocale *lie-group charts tms one mns invs*
using *is-lie-group* .

end

lemma *lie-group-imp-lie-grp:*
assumes *lie-group charts pls one any-mns any-invs*
shows *lie-grp charts pls one*
unfolding *lie-grp-def lie-grp-axioms-def* **apply** (*intro conjI*)
subgoal using *assms lie-group-def* **by** *blast*
subgoal
using *assms* **unfolding** *grp-on-def grp-on-axioms-def lie-group-def group-on-with-def group-on-with-axioms-def*
by (*meson assms group-on-with.right-minus lie-group.axioms(2)*)
subgoal using *assms* **unfolding** *lie-group-def lie-group-axioms-def* **by** *simp*

```

subgoal using assms unfolding lie-group-def lie-group-axioms-def
  by (smt (verit, ccfv-threshold) ⟨grp-on (manifold.carrier charts) pls one⟩
diff.diff-cong
  group-on-with.inv-is-unique group-on-with.right-minus group-on-with.uminus-mem
grp-on.is-group-on-with)
done

```

We give a few intro rules for the *lie-group* predicate, as well as an Eisbach method for further breaking down the proof of smoothness of the multiplication and inversion maps. This should lead to fairly organised proofs that some structure is a *lie-group*. In general, I would prefer *group-manifold-imp-lie-group2* to *group-manifold-imp-lie-group*.

```

lemma group-manifold-imp-lie-group [intro]:
  assumes is-manifold: c-manifold c ∞
    and is-group: group-on-with (⋃ (domain ‘ c)) tms tms-1 dvsn invs
    and smooth-mult: diff ∞ (c-manifold-prod.prod-charts c c) c (λ(a,b). tms a b)
    and smooth-inv: diff ∞ c c invs
  shows lie-group c tms tms-1 dvsn invs
  unfolding lie-group-def manifold.carrier-def lie-group-axioms-def
  by (simp-all add: c-manifold-prod-def is-manifold is-group smooth-inv smooth-mult)

```

```

lemma group-manifold-imp-lie-group2 [intro]:
  assumes is-manifold: c-manifold c ∞
    and is-group: group-on-with (⋃ (domain ‘ c)) tms tms-1 dvsn invs
    and smooth-mult: diff-axioms ∞ (c-manifold-prod.prod-charts c c) c (λ(a,b). tms a b)
    and smooth-inv: diff-axioms ∞ c c invs
  shows lie-group c tms tms-1 dvsn invs
  by (auto intro!: c-manifolds.intro diff.intro simp: assms c-manifold-prod.c-manifold-atlas-product c-manifold-prod-def)

```

```

lemma lie-grpI [intro]:
  fixes tms tms-1 c
  defines invs ≡ grp-on.invs (⋃ (domain ‘ c)) tms tms-1
  assumes is-manifold: c-manifold c ∞
    and is-group: grp-on (⋃ (domain ‘ c)) tms tms-1
    and smooth-mult: diff-axioms ∞ (c-manifold-prod.prod-charts c c) c (λ(a,b). tms a b)
    and smooth-inv: diff-axioms ∞ c c invs
  shows lie-grp c tms tms-1
  by (metis group-manifold-imp-lie-group2 grp-on.is-group-on-with invs-def is-group is-manifold lie-group-imp-lie-grp smooth-inv smooth-mult)

```

A small method to unfold the axioms of differentiability of group operations. Allows for succinct goals to be stated while quickly unfolding to a useful level of technicality.

```

method unfold-diff-axioms = (
  unfold diff-axioms-def,

```

```

    rule allI,
    rule impI,
    (rule bezI)+,
    (rule conjI),
    rule-tac[2] conjI
  )

```

7.3 Some lemmas about Lie groups (and other needed results).

context *lie-group* **begin**

lemma *obtain-chart-cover*:

assumes $S \subseteq \text{carrier}$

obtains C **where** $\forall c \in C. c \in \text{atlas} \ \forall s \in S. \exists c \in C. s \in \text{domain } c$

by (*metis assms carrierE in-charts-in-atlas subset-iff*)

lemma *open-covered-by-charts*:

assumes $S \subseteq \text{carrier}$ *open* S

obtains C **where** $\forall c \in C. c \in \text{atlas} \ S = \bigcup \{ \text{domain } c \mid c. c \in C \}$

proof –

obtain C **where** $C: \forall c \in C. c \in \text{atlas} \ \forall s \in S. \exists c \in C. s \in \text{domain } c$

using *obtain-chart-cover assms* **by** *blast*

let $?restr\text{-}chart = \lambda c. \text{if } \text{domain } c \subseteq S \text{ then } c \text{ else } \text{restrict-chart } S \ c$

let $?C = \{ ?restr\text{-}chart \ c \mid c. c \in C \}$

have $\forall c \in ?C. c \in \text{atlas}$

using $C(1)$ *restrict-chart-in-atlas* **by** *auto*

moreover **have** $S = \bigcup \{ \text{domain } c \mid c. c \in ?C \}$

using $\text{assms}(2)$ *domain-restrict-chart* **by** (*auto, metis C(2) Int-iff, fastforce*)

ultimately show *?thesis* **using** *that* **by** *presburger*

qed

lemma *lie-prod: c-manifold-prod* ∞ *charts charts*

by *unfold-locales*

interpretation *lie-prod: c-manifold-prod* ∞ *charts charts*

by *unfold-locales*

lemma *continuous-on-tms*:

assumes $x \in \text{carrier}$

shows *continuous-on carrier* $(\lambda y. \text{tms } x \ y)$

and *continuous-on carrier* $(\lambda y. \text{tms } y \ x)$

proof –

have *cts-tms: continuous-on lie-prod.carrier* $(\lambda(a, b). \text{tms } a \ b)$

using *lie-group-axioms diff.continuous-on unfolding lie-group-def lie-group-axioms-def*

by *blast*

have *tms-is-comp*: $(\text{tms } x) = (\lambda(a, b). \text{tms } a \ b) \circ (\lambda y. (x, y))$

by (*simp add: comp-def*)

show *continuous-on carrier* $(\lambda y. \text{tms } x \ y)$

```

proof –
  have cts-R: continuous-on carrier ( $\lambda y. (x,y)$ )
    using continuous-on-Pair[OF continuous-on-const[of carrier x] continuous-on-id] .
  have pair-carrier: Pair x ‘ carrier  $\subseteq$  lie-prod.carrier
    unfolding image-def using lie-prod.prod-carrier assms by blast
  thus ?thesis
  using continuous-on-compose[OF cts-R] cts-tms tms-is-comp continuous-on-subset[OF
- pair-carrier]
    by metis
  qed
  show continuous-on carrier ( $\lambda y. tms\ y\ x$ )
proof –
  have cts-L: continuous-on carrier ( $\lambda y. (y,x)$ )
    using continuous-on-Pair[OF continuous-on-id continuous-on-const[of carrier
x]] .
  have pair-carrier': ( $\lambda y. (y,x)$ ) ‘ carrier  $\subseteq$  lie-prod.carrier
    unfolding image-def using lie-prod.prod-carrier assms by blast
  thus ?thesis
  using continuous-on-compose[OF cts-L] cts-tms tms-is-comp continuous-on-subset[OF
- pair-carrier']
    by force
  qed
qed

```

```

lemma diff-tms:
  assumes  $x \in carrier$ 
  shows diff  $\infty$  charts charts ( $\lambda y. tms\ x\ y$ )
    and diff  $\infty$  charts charts ( $\lambda y. tms\ y\ x$ )
  subgoal
    using diff-compose[OF lie-prod.diff-left-Pair[OF assms] smooth-mult] diff.diff-cong
by fastforce
  subgoal
    using diff-compose[OF lie-prod.diff-right-Pair[OF assms] smooth-mult] diff.diff-cong
by fastforce
  done

```

```

lemma diff-tms-invs:
  assumes  $x \in carrier$ 
  shows diff  $\infty$  charts charts ( $\lambda y. tms\ (invs\ x)\ y$ )
    and diff  $\infty$  charts charts ( $\lambda y. tms\ y\ (invs\ x)$ )
  using diff-tms[of invs x] assms uminus-mem by blast+

```

```

lemma diff-tms-invs':
  assumes  $x \in carrier$ 
  shows diff  $\infty$  charts charts ( $\lambda y. tms\ x\ (invs\ y)$ )
    and diff  $\infty$  charts charts ( $\lambda y. tms\ (invs\ y)\ x$ )
  using diff-compose[OF smooth-inv diff-tms(1)[OF assms]] apply (simp add:
diff.diff-cong)

```


using *diff-compose*[*OF smooth-inv diff-tms*(2)[*OF assms*]] **by** (*simp add: diff.diff-cong*)

end

8 Morphisms of Lie groups, actions and representations

8.1 Morphism of Lie groups.

locale *lie-group-pair* =

L1: lie-group c1 t1 i1 d1 m1 +

L2: lie-group c2 t2 i2 d2 m2

for *c1* :: ('a::{\i>second-countable-topology,t2-space}, 'b::\i>euclidean-space) *chart set*

and *c2* :: ('c::{\i>second-countable-topology,t2-space}, 'd::\i>euclidean-space) *chart set*

and *t1 t2 and i1 i2 and d1 d2 and m1 m2*

locale *lie-group-morphism-with* =

lie-group-pair c1 c2 t1 t2 i1 i2 d1 d2 m1 m2 +

diff ∞ *c1 c2 f* +

group-hom-betw L1.carrier L2.carrier t1 t2 i1 i2 d1 d2 m1 m2 f

for *c1* :: ('a::{\i>second-countable-topology,t2-space}, 'b::\i>euclidean-space) *chart set*

and *c2* :: ('c::{\i>second-countable-topology,t2-space}, 'd::\i>euclidean-space) *chart set*

and *t1 t2 and i1 i2 and d1 d2 and m1 m2 and f*

lemma (**in** *lie-group-pair*) *lie-group-morphismI*:

assumes *diff* ∞ *c1 c2 f*

and *group-hom*: $\forall x \in L1.carrier. \forall y \in L1.carrier. f (t1 x y) = t2 (f x) (f y)$

and *closure*: $\forall x \in L1.carrier. f x \in L2.carrier$

shows *lie-group-morphism-with c1 c2 t1 t2 i1 i2 d1 d2 m1 m2 f*

proof –

have *1: group-on-with-pair L1.carrier L2.carrier t1 t2 i1 i2 d1 d2 m1 m2*

using *lie-group-pair-axioms unfolding lie-group-pair-def lie-group-def group-on-with-pair-def*

by *presburger*

show *?thesis*

unfolding *lie-group-morphism-with-def group-hom-betw-def group-hom-betw-axioms-def*

by (*simp add: assms lie-group-pair-axioms 1*)

qed

lemma (**in** *lie-group*) *lie-group-morphismI*:

assumes *lie-group c2 t2 i2 d2 m2*

and *diff* ∞ *charts c2 f*

and *group-hom*: $\forall x \in carrier. \forall y \in carrier. f (tms x y) = t2 (f x) (f y)$

and *closure*: $\forall x \in carrier. f x \in (manifold.carrier c2)$

shows *lie-group-morphism-with charts c2 tms t2 tms-one i2 d2 m2 f*

by (*auto intro: lie-group-pair.lie-group-morphismI simp: lie-group-pair-def lie-group-axioms assms*)

```

locale lie-group-isomorphism =
  lie-group-pair c1 c2 t1 t2 i1 i2 d1 d2 m1 m2 +
  diffeomorphism  $\infty$  c1 c2 f f' +
  group-hom-betw L1.carrier L2.carrier t1 t2 i1 i2 d1 d2 m1 m2 f
  for c1 :: ('a::{second-countable-topology,t2-space}, 'b::euclidean-space) chart set
    and c2 :: ('c::{second-countable-topology,t2-space}, 'd::euclidean-space) chart set
  and t1 t2 and i1 i2 and d1 d2 and m1 m2 and f f'

```

8.2 Action of a Lie group on a manifold.

```

abbreviation (input) diff-action-map g-charts m-charts action  $\equiv$ 
  diff  $\infty$  (c-manifold-prod.prod-charts g-charts m-charts) m-charts action

```

A Lie group action is a homomorphism from the Lie group to the automorphism group of a space, here a manifold, which is differentiable (smooth). I take here the more explicit definition given in Kirillov's lecture notes (2008; page 12), and derive the more abstract version later (after showing *c-manifold.Diff* is not just a group, but a Lie group).

Take care: there are now two manifolds, of which the Lie group is the primary one as far as namespace is concerned. Everything pertaining to the manifold acted upon is accessed with qualified syntax. This disappears for Lie groups acting on themselves.

```

locale lie-group-action =
  lie-group charts tms tms-one dvsn invs + M: c-manifold m-charts k
  for charts::('a::{t2-space,second-countable-topology}, 'e::euclidean-space) chart set
    and tms tms-one dvsn invs
    and m-charts::('b::{t2-space,second-countable-topology}, 'f::euclidean-space) chart set and k +
  fixes action ( $\langle \varrho \rangle$ )
  assumes act-diff:  $g \in \text{carrier} \implies (\varrho g) \in M.\text{Diff}$ 
    and act-one:  $\varrho \text{ tms-one} = M.\text{Diff-id}$ 
    and act-hom:  $f \in G \implies g \in G \implies \varrho (\text{tms } f g) = M.\text{Diff-comp } (\varrho f) (\varrho g)$ 
    and act-diff-prod: diff-action-map charts m-charts ( $\lambda(g,m). \text{the } ((\varrho g) m)$ )

```

After proving Diff is a group, some of these axioms can be replaced.

```

locale lie-group-action' =
  lie-group charts tms tms-one dvsn invs +
  M: c-manifold m-charts k +
  A: group-hom-betw carrier M.Diff tms M.Diff-comp tms-one M.Diff-id dvsn
  M.Diff-comp-inv invs M.Diff-inv  $\varrho$ 
  for charts::('a::{t2-space,second-countable-topology}, 'e::euclidean-space) chart set
    and tms tms-one dvsn invs
    and m-charts::('b::{t2-space,second-countable-topology}, 'f::euclidean-space) chart set and k
  and  $\varrho$  :: 'a  $\Rightarrow$  ('b  $\rightarrow$  'b) +

```

assumes *diff-action-map*: *diff-action-map charts m-charts* ($\lambda(g,m)$. the $((\rho g m))$)

8.3 Action of a Lie Group on itself.

context *lie-group begin*

abbreviation (*input*) *left-self-action* :: $'a \Rightarrow 'a \Rightarrow 'a$ ($\langle \mathcal{L} \rightarrow [91]$)
where *left-self-action* $g g' \equiv \text{tms } g g'$

abbreviation *left-action* :: $'a \Rightarrow ('a \rightarrow 'a)$
where *left-action* $g \equiv (\lambda x$. if $x \in \text{carrier}$ then *Some* (*left-self-action* $g x$) else *None*)

abbreviation (*input*) *right-self-action* :: $'a \Rightarrow 'a \Rightarrow 'a$ ($\langle \mathcal{R} \rightarrow [91]$)
where *right-self-action* $g g' \equiv \text{tms } g' (\text{invs } g)$

abbreviation *right-action* :: $'a \Rightarrow ('a \rightarrow 'a)$
where *right-action* $g \equiv (\lambda x$. if $x \in \text{carrier}$ then *Some* (*right-self-action* $g x$) else *None*)

abbreviation (*input*) *adjoint-self-action* :: $'a \Rightarrow 'a \Rightarrow 'a$
where *adjoint-self-action* $g g' \equiv \text{tms } g (\text{tms } g' (\text{invs } g))$

8.3.1 The left action.

lemma *L-action-in*: $(\text{left-self-action } g g') \in \text{carrier}$ **if** $g \in \text{carrier } g' \in \text{carrier}$
by (*simp add: add-mem that*)

lemma *the-left-action*: *left-self-action* $x y = \text{the } (\text{left-action } x y)$ **if** $y \in \text{carrier}$
by (*simp add: that*)

lemma *L-action-invs*: $(\text{left-self-action } (\text{invs } x) \circ \text{left-self-action } x) y = y$
 $(\text{left-self-action } x \circ \text{left-self-action } (\text{invs } x)) y = y$
if $x \in \text{carrier } y \in \text{carrier}$
apply (*metis (no-types, lifting) add-assoc add-zeroL comp-apply left-minus that uminus-mem*)
by (*metis (no-types, lifting) add-assoc add-zeroL comp-apply right-minus that uminus-mem*)

lemma *L-homeomorphism*: *homeomorphism carrier carrier* ($\mathcal{L} x$) ($\mathcal{L} (\text{invs } x)$) **if** $x \in \text{carrier}$

proof –

{
fix $x y$ **assume** *xy-in-carrier*: $x \in \text{carrier } y \in \text{carrier}$
then have $\text{tms } (\text{invs } x) (\text{tms } x y) = y$ **and** $\text{tms } x (\text{tms } (\text{invs } x) y) = y$
using *add-assoc add-zeroL uminus-mem* **by** (*metis left-minus, metis right-minus*)
}
thus *homeomorphism carrier carrier* ($\text{tms } x$) ($\text{tms } (\text{invs } x)$)
using *that continuous-on-tms(1)* **by** (*auto intro: homeomorphismI simp: L-action-in image-subset-iff uminus-mem*)

qed

lemma *L-homeomorphism'*: homeomorphism carrier carrier $(\mathcal{L} (invs x)) (\mathcal{L} x)$
if $x \in carrier$
using *L-homeomorphism homeomorphism-sym that* **by** blast

lemma *L-homeomorphism-chart*: homeomorphism (domain c) $(\mathcal{L} x \text{ ' } domain c) (\mathcal{L} x)$
if $x \in carrier c \in atlas$
using *L-homeomorphism homeomorphism-of-subsets that* **by** blast

lemma *L-homeomorphism-chart'*: homeomorphism $(\mathcal{L} x \text{ ' } domain c) (domain c)$
 $(\mathcal{L} (invs x)) (\mathcal{L} x)$
if $x \in carrier c \in atlas$
using *L-homeomorphism-chart that homeomorphism-sym* **by** blast

lemma *L-open-map*:
assumes $x \in carrier$ open S $S \subseteq carrier$
shows open $(\mathcal{L} x \text{ ' } S)$
proof –
obtain C **where** C: $\forall c \in C. c \in atlas S = \bigcup \{ domain c \mid c. c \in C \}$
using *open-covered-by-charts assms* **by** blast
have $\mathcal{L} x \text{ ' } S = \bigcup \{ \mathcal{L} x \text{ ' } domain c \mid c. c \in C \}$
using C(2) **by** auto
thus open $(\mathcal{L} x \text{ ' } S)$
using *homeomorphism-imp-open-map' L-homeomorphism* **by** (metis assms open-carrier)
qed

lift-definition *L-chart* :: 'a \Rightarrow ('a,'e) chart \Rightarrow ('a,'e) chart
is $\lambda x. \lambda (d,d',f,f').$ if $x \in carrier \wedge d \subseteq carrier$ then $(\mathcal{L} x \text{ ' } d, d', f \circ \mathcal{L} (invs x), \mathcal{L} x \circ f')$ else $(\{\}, \{\}, f, f')$
using *L-homeomorphism* **by** (auto split: if-splits intro!: L-open-map)
(meson homeomorphism-compose homeomorphism-of-subsets homeomorphism-symD)

lemma *L-chart-apply-chart[simp]*: apply-chart (L-chart x c) = apply-chart c $\circ \mathcal{L} (invs x)$
and *L-chart-inv-chart[simp]*: inv-chart (L-chart x c) = $\mathcal{L} x \circ inv-chart c$
and *domain-L-chart[simp]*: domain (L-chart x c) = $\mathcal{L} x \text{ ' } domain c$
and *codomain-L-chart[simp]*: codomain (L-chart x c) = codomain c
if $x \in carrier c \in atlas$
using *that(1) domain-atlas-subset-carrier[OF that(2)]* **by** (transfer, auto)+

lemma *L-chart-apply-chart'[simp]*: apply-chart (L-chart x c) = apply-chart c $\circ \mathcal{L} (invs x)$
and *L-chart-inv-chart'[simp]*: inv-chart (L-chart x c) = $\mathcal{L} x \circ inv-chart c$
and *domain-L-chart'[simp]*: domain (L-chart x c) = $\mathcal{L} x \text{ ' } domain c$
and *codomain-L-chart'[simp]*: codomain (L-chart x c) = codomain c
if $x \in carrier domain c \subseteq carrier$

```

using that by (transfer, auto)+

lemma smooth-compat-L-chart:
  assumes  $x \in \text{carrier } c \in \text{atlas } c' \in \text{atlas}$ 
  shows  $\infty\text{-smooth-compat } (L\text{-chart } x \ c) \ c'$ 
proof -
  let ?dom1 =  $(\lambda y. c \ (tms \ (invs \ x) \ y)) \ ' \ (tms \ x \ ' \ \text{domain } c \ \cap \ \text{domain } c')$ 
  let ?dom2 =  $\text{codomain } c \ \cap \ \text{inv-chart } c \ - \ ' \ (\text{carrier} \ \cap \ \text{tms } x \ - \ ' \ \text{domain } c')$ 
  let ?dom3 =  $c' \ ' \ (tms \ x \ ' \ \text{domain } c \ \cap \ \text{domain } c')$ 
  let ?dom4 =  $\text{codomain } c' \ \cap \ \text{inv-chart } c' \ - \ ' \ (\text{carrier} \ \cap \ \text{tms} \ (invs \ x) \ - \ ' \ \text{domain } c)$ 

  have invs-tms-defined:  $c \ (tms \ (invs \ x) \ (tms \ x \ y)) \ \in \ \text{codomain } c$  if  $y \in \text{domain } c$ 
for  $y$ 
  by (metis add-assoc add-uminus add-zeroL assms(1,2) chart-in-codomain in-carrier-atlasI uminus-mem that)
  have domain-simp-1: ?dom1 = ?dom2
proof -
  {
    fix  $y$  assume  $y: tms \ x \ y \ \in \ \text{domain } c' \ y \ \in \ \text{domain } c$ 
    have inv-chart c  $(c \ (tms \ (invs \ x) \ (tms \ x \ y))) \ \in \ \text{carrier}$ 
      and  $tms \ x \ (inv\text{-chart } c \ (c \ (tms \ (invs \ x) \ (tms \ x \ y)))) \ \in \ \text{domain } c'$ 
    subgoal using  $y$  assms(2) invs-tms-defined by blast
    subgoal using  $y$  by (metis add-assoc add-zeroL assms(1,2) in-carrier-atlasI inv-chart-inverse left-minus uminus-mem)
    done
  } moreover {
    fix  $y$  assume  $y: y \ \in \ \text{codomain } c \ \text{inv-chart } c \ y \ \in \ \text{carrier } tms \ x \ (inv\text{-chart } c \ y)$ 
 $\in \ \text{domain } c'$ 
    have  $y = c \ (tms \ (invs \ x) \ (tms \ x \ (inv\text{-chart } c \ y)))$ 
    by (metis (full-types) assms(1) chart-inverse-inv-chart homeomorphism-apply1 L-homeomorphism y(1,2))
    then have  $y \ \in \ (\lambda y. c \ (tms \ (invs \ x) \ y)) \ ' \ (tms \ x \ ' \ \text{domain } c \ \cap \ \text{domain } c')$ 
      using  $y(1,3)$  by blast
  }
  ultimately show ?dom1 = ?dom2 using invs-tms-defined by auto
qed
  have domain-simp-2: ?dom3 = ?dom4
proof -
  {
    fix  $y$  assume  $y: tms \ x \ y \ \in \ \text{domain } c' \ y \ \in \ \text{domain } c$ 
    have  $tms \ x \ y \ \in \ \text{carrier}$  and  $tms \ (invs \ x) \ (tms \ x \ y) \ \in \ \text{domain } c$ 
    subgoal using  $y$  assms(3) by simp
    subgoal using  $y$  by (metis add-assoc add-zeroL assms(1,2) in-carrier-atlasI local.left-minus uminus-mem)
    done
  } moreover {
    fix  $y$  assume  $y \ \in \ \text{codomain } c' \ \text{inv-chart } c' \ y \ \in \ \text{carrier } tms \ (invs \ x) \ (inv\text{-chart } c' \ y) \ \in \ \text{domain } c$ 
    then have  $y \ \in \ c' \ ' \ (tms \ x \ ' \ \text{domain } c \ \cap \ \text{domain } c')$ 
  }

```

by (*smt* (*verit*, *ccfv-threshold*) *Int-iff add-assoc add-uminus add-zeroL*
assms(1)
chart-inverse-inv-chart inv-chart-in-domain rev-image-eqI uminus-mem)
}
ultimately show $?dom3 = ?dom4$ **by** *auto*
qed

have *smooth-on* $?dom1$ ($c' \circ (tms\ x \circ inv\text{-}chart\ c)$)
using *diff.diff-chartsD*[*OF diff-tms(1)*][*OF assms(1)*] *assms(2,3)*
by (*simp add: comp-assoc domain-simp-1*)
moreover have *smooth-on* $?dom3$ ($c \circ tms\ (invs\ x) \circ inv\text{-}chart\ c'$)
using *diff.diff-chartsD*[*OF diff-tms-invs(1)*][*OF assms(1)*] *assms(3,2)* **by** (*simp*
add: domain-simp-2)
ultimately show *?thesis*
by (*unfold smooth-compat-def, auto simp: assms*)
qed

lemma *L-chart-compat*:
assumes $x \in carrier\ c \in atlas$
shows $\infty\text{-smooth-compat}\ c\ (L\text{-chart}\ x\ c)$
using *smooth-compat-L-chart*[*OF assms(1,2,2)*] **by** (*simp add: smooth-compat-commute*)

lemma *L-chart-in-atlas*: $L\text{-chart}\ x\ c \in atlas$ **if** $x \in carrier\ c \in atlas$
proof (*rule maximal-atlas*)
show $domain\ (L\text{-chart}\ x\ c) \subseteq carrier$ **using** *L-action-in* **that** **by** *auto*
fix c' **assume** $c' \in atlas$
with *that(2)* **have** $\infty\text{-smooth-compat}\ c\ c'$ **by** (*simp add: atlas-is-atlas*)
thus $\infty\text{-smooth-compat}\ (L\text{-chart}\ x\ c)\ c'$
using *smooth-compat-L-chart*[*OF that*] **by** (*simp add: <c' ∈ atlas>*)
qed

lemma *left-action-automorphic*: $c\text{-automorphism}\ \infty\ charts\ (\mathcal{L}\ x)\ (\mathcal{L}\ (invs\ x))$
if $x \in carrier$
proof (*unfold-locales*)
fix y **assume** $y \in carrier$
then obtain $c1$ **where** $c1: c1 \in atlas\ y \in domain\ c1$ **using** *atlasE* **by** *blast*
let $?L = left\text{-self-action}\ x$
let $?L_i = left\text{-self-action}\ (invs\ x)$

To find the second chart, for the codomain of $tms\ x$, just shift the first chart across.

show $\exists c1 \in atlas. \exists c2 \in atlas.$
 $y \in domain\ c1 \wedge$
 $?L\ ' domain\ c1 \subseteq domain\ c2 \wedge$
 $smooth\text{-on}\ (codomain\ c1)\ (c2 \circ ?L \circ inv\text{-}chart\ c1)$
proof (*intro bexI conjI*)
let $?c2 = L\text{-chart}\ x\ c1$

show $y \in domain\ c1$ **by** (*simp add: c1(2)*)

show $c1 \in \text{atlas } ?c2 \in \text{atlas}$ **by** (*simp add: L-chart-in-atlas c1(1) that*)+
show $\text{tms } x \text{ ' domain } c1 \subseteq \text{domain } ?c2$ **by** (*simp add: c1(1) that*)

have $(c1 \circ ?L_i \circ ?L \circ \text{inv-chart } c1) a = a$ **if** $a \in \text{codomain } c1$ **for** a
using *L-action-invs(1) <x ∈ carrier> c1(1) that by force*
thus *smooth-on (codomain c1) (?c2 ∘ ?L ∘ inv-chart c1)*
using *smooth-on-id smooth-on-cong*
by (*smt (verit, del-insts) L-chart-apply-chart c1(1) open-codomain that*)
qed

show $\exists c1 \in \text{atlas}. \exists c2 \in \text{atlas}.$

$y \in \text{domain } c1 \wedge$
 $?L_i \text{ ' domain } c1 \subseteq \text{domain } c2 \wedge$
 $\text{smooth-on (codomain } c1) (c2 \circ ?L_i \circ \text{inv-chart } c1)$

proof (*intro bexI conjI*)

let $?c2 = \text{L-chart (invs } x) c1$

have [*simp*]: $\text{invs } x \in \text{carrier}$ **by** (*simp add: that uminus-mem*)

show $y \in \text{domain } c1$ **by** (*simp add: c1(2)*)

show $c1 \in \text{atlas } ?c2 \in \text{atlas}$ **by** (*simp add: L-chart-in-atlas c1(1) that*)+

show $\text{tms (invs } x) \text{ ' domain } c1 \subseteq \text{domain } ?c2$ **by** (*simp add: c1(1) that*)

have $1: (c1 \circ ?L \circ ?L_i \circ \text{inv-chart } c1) a = a$

if $a \in \text{codomain } c1$ **for** a

using *L-action-invs(2) <x ∈ carrier> c1 that by force*

show *smooth-on (codomain c1) (?c2 ∘ ?L_i ∘ inv-chart c1)*

apply (*simp add: c1 uminus-uminus[OF that]*)

using *smooth-on-id 1 by (smt (verit, del-insts) open-codomain smooth-on-cong)*

qed

{ **fix** y **assume** $y \in \text{carrier}$

show $\text{tms (invs } x) (\text{tms } x y) = y$

by (*metis <y ∈ carrier> add-assoc add-zeroL left-minus that uminus-mem*)

show $\text{tms } x (\text{tms (invs } x) y) = y$

by (*metis <y ∈ carrier> add-assoc add-zeroL right-minus that uminus-mem*) }

qed

lemma *left-action-in-Diff: left-action* $x \in \text{Diff}$ **if** $x \in \text{carrier}$

apply (*intro DiffI automorphismI exI[where x=left-self-action (invs x)]*)

subgoal using *c-automorphism.c-automorphism-cong left-action-automorphic that by fastforce*

subgoal by (*simp add: domIff order-class.order-eq-iff subset-iff*)

done

lemma *diff-the-L: diff* $\infty (c\text{-manifold-prod.prod-charts charts charts}) \text{charts } (\lambda(g, m). \text{the (left-action } g m))$

(*is diff* $\infty ?\text{prod-charts charts } ?L$)

proof –

```

let ?prod-carrier = manifold.carrier ?prod-charts
have L-eq: ?L (g,m) = (L g) m if (g,m) ∈ ?prod-carrier for g m
  using c-manifold-prod.prod-carrier[OF lie-prod] that by fastforce
show ?thesis
  apply (rule diff.diff-cong[OF smooth-mult])
  using L-eq by fastforce
qed

```

```

lemma left-action: lie-group-action' charts tms tms-one dvn invs charts ∞ left-action
  unfolding lie-group-action'-def lie-group-action'-axioms-def
  apply (simp add: lie-group-axioms c-manifold-axioms, intro conjI)
  subgoal using add-assoc add-mem left-action-in-Diff by (unfold-locales, auto)
  subgoal by (rule diff-the-L)
done

```

```

sublocale left-action: lie-group-action' charts tms tms-one dvn invs charts ∞
left-action
  by (rule left-action)

```

8.3.2 The right action.

```

lemma R-action-in: (right-self-action g g') ∈ carrier if g ∈ carrier g' ∈ carrier
  by (simp add: add-mem that uminus-mem)

```

```

lemma the-right-action: right-self-action x y = the (right-action x y) if y ∈ carrier
  by (simp add: that)

```

```

lemma R-action-invs: (right-self-action (invs x) ∘ right-self-action x) y = y
  (right-self-action x ∘ right-self-action (invs x)) y = y
  if x ∈ carrier y ∈ carrier
  using add-assoc add-zeroR comp-apply right-minus left-minus that uminus-mem
  by simp-all

```

```

lemma R-homeomorphism: homeomorphism carrier carrier (R x) (R (invs x))
  if x ∈ carrier

```

```

proof -
  {
    fix x y assume xy-in-carrier: x ∈ carrier y ∈ carrier
    then have tms (tms y (invs x)) (invs (invs x)) = y and tms (tms y (invs (invs
x))) (invs x) = y
    using add-assoc add-zeroR uminus-mem by (metis right-minus, metis left-minus)
  }
  thus homeomorphism carrier carrier (R x) (R (invs x))
    using that continuous-on-tms(2) by (auto intro!: homeomorphismI simp:
R-action-in image-subset-iff uminus-mem)
qed

```

```

lemma R-homeomorphism': homeomorphism carrier carrier (R (invs x)) (R x)
  if x ∈ carrier

```


using *R-homeomorphism homeomorphism-sym that by blast*

lemma *R-homeomorphism-chart: homeomorphism (domain c) (R x ' domain c)*
(R x) (R (invs x))
if $x \in \text{carrier } c \in \text{atlas}$
using *R-homeomorphism homeomorphism-of-subsets that by blast*

lemma *R-homeomorphism-chart': homeomorphism (R x ' domain c) (domain c)*
(R (invs x)) (R x)
if $x \in \text{carrier } c \in \text{atlas}$
using *R-homeomorphism-chart that homeomorphism-sym by blast*

lemma *R-open-map:*

assumes $x \in \text{carrier } \text{open } S \ S \subseteq \text{carrier}$

shows *open (R x ' S)*

proof –

obtain *C where C: $\forall c \in C. c \in \text{atlas } S = \bigcup \{\text{domain } c \mid c. c \in C\}$*

using *open-covered-by-charts assms by blast*

have $\mathcal{R} \ x \ ' \ S = \bigcup \{\mathcal{R} \ x \ ' \ \text{domain } c \mid c. c \in C\}$

using *C(2) by auto*

thus *open (R x ' S)*

using *homeomorphism-imp-open-map' R-homeomorphism assms open-carrier*

by *fast*

qed

lift-definition *R-chart :: 'a \Rightarrow ('a,'e) chart \Rightarrow ('a,'e) chart*

is $\lambda x. \lambda(d,d',f,f'). \text{if } x \in \text{carrier} \wedge d \subseteq \text{carrier} \text{ then } (\mathcal{R} \ x \ ' \ d, d', f \circ \mathcal{R} \ (\text{invs } x), \mathcal{R} \ x \circ f') \text{ else } (\{\}, \{\}, f, f')$

using *R-homeomorphism by (auto split: if-splits intro!: R-open-map)*

(meson homeomorphism-compose homeomorphism-of-subsets homeomorphism-symD)

lemma *R-chart-apply-chart[simp]: apply-chart (R-chart x c) = apply-chart c \circ R*
(invs x)

and *R-chart-inv-chart[simp]: inv-chart (R-chart x c) = R x \circ inv-chart c*

and *domain-R-chart[simp]: domain (R-chart x c) = R x ' domain c*

and *codomain-R-chart[simp]: codomain (R-chart x c) = codomain c*

if $x \in \text{carrier } c \in \text{atlas}$

using *that(1) domain-atlas-subset-carrier[OF that(2)] by (transfer, auto)+*

lemma *R-chart-apply-chart'[simp]: apply-chart (R-chart x c) = apply-chart c \circ R*
(invs x)

and *R-chart-inv-chart'[simp]: inv-chart (R-chart x c) = R x \circ inv-chart c*

and *domain-R-chart'[simp]: domain (R-chart x c) = R x ' domain c*

and *codomain-R-chart'[simp]: codomain (R-chart x c) = codomain c*

if $x \in \text{carrier } \text{domain } c \subseteq \text{carrier}$

using *that by (transfer, auto)+*

lemma *smooth-compat-R-chart:*

assumes $x \in \text{carrier } c \in \text{atlas } c' \in \text{atlas}$

shows ∞ -smooth-compat (R-chart x c) c'

proof –

let ?dom1 = ($\lambda y. c (tms\ y (invs\ (invs\ x)))$) ‘ ($\lambda g'. tms\ g' (invs\ x)$) ‘ domain c
 \cap domain c')

let ?dom2 = codomain c \cap inv-chart c – ‘ (carrier \cap ($\lambda y. tms\ y (invs\ x)$) – ‘
 domain c')

let ?dom3 = c' ‘ ($\lambda y. tms\ y (invs\ x)$) ‘ domain c \cap domain c')

let ?dom4 = codomain c' \cap inv-chart c' – ‘ (carrier \cap ($\lambda y. tms\ y\ x$) – ‘ domain
 c)

have invs-tms-defined: c (tms (tms y (invs x)) (invs (invs x))) \in codomain c **if**
 y \in domain c **for** y

using add-assoc add-zeroR assms(1,2) local.right-minus that uminus-mem **by**
 auto

then have domain-simp-1: ?dom1 = ?dom2

proof –

{

fix y **assume** y: tms y (invs x) \in domain c' y \in domain c

have inv-chart c (c (tms (tms y (invs x)) (invs (invs x)))) \in carrier

and tms (inv-chart c (c (tms (tms y (invs x)) (invs (invs x))))) (invs x) \in
 domain c'

subgoal using y invs-tms-defined assms(2) **by** blast

subgoal using y add-assoc add-zeroR assms(1,2) in-carrier-atlasI inv-chart-inverse
 right-minus uminus-mem **by** metis

done

} **moreover** {

fix y **assume** y \in codomain c inv-chart c y \in carrier tms (inv-chart c y) (invs
 x) \in domain c'

then have y \in ($\lambda y. apply-chart\ c (tms\ y (invs\ (invs\ x)))$) ‘ ($\lambda g'. tms\ g' (invs$
 x)) ‘ domain c \cap domain c')

by (smt (verit, cfv-threshold) IntI R-action-invs(1) assms(1) chart-inverse-inv-chart
 comp-apply imageI inv-chart-in-domain)

}

ultimately show ?dom1 = ?dom2 **using** invs-tms-defined **by** auto

qed

have domain-simp-2: ?dom3 = ?dom4

proof –

{

fix y **assume** y: tms y (invs x) \in domain c' y \in domain c

have tms y (invs x) \in carrier **and** tms (tms y (invs x)) x \in domain c

subgoal using y assms(3) **by** simp

subgoal using y **by** (metis add-assoc add-zeroR assms(1,2) in-carrier-atlasI
 left-minus uminus-mem)

done

} **moreover** {

fix xa **assume** xa: xa \in codomain c' inv-chart c' xa \in carrier tms (inv-chart
 c' xa) x \in domain c

then have xa \in apply-chart c' ‘ ($\lambda y. tms\ y (invs\ x)$) ‘ domain c \cap domain
 c')

```

    by (smt (verit, ccfv-threshold) Int-iff add-assoc add-zero assms(1) chart-inverse-inv-chart
        inv-chart-in-domain right-minus rev-image-eqI uminus-mem)
  }
  ultimately show ?dom3 = ?dom4 by auto
qed

```

```

have smooth-on ?dom1 (c' ∘ ((λg'. tms g' (invs x)) ∘ inv-chart c))
  using diff.diff-chartsD[OF diff-tms-invs(2)[OF assms(1)] assms(2,3)]
  by (simp add: comp-assoc domain-simp-1)
moreover have smooth-on ?dom3 (c ∘ (λg'. tms g' (invs (invs x)))) ∘ inv-chart
c')
  using diff.diff-chartsD[OF diff-tms(2)] uminus-uminus assms by (simp add:
domain-simp-2)
  ultimately show ?thesis
  by (unfold smooth-compat-def, auto simp: assms)
qed

```

```

lemma R-chart-compat:
  assumes x ∈ carrier c ∈ atlas
  shows ∞-smooth-compat c (R-chart x c)
  using smooth-compat-R-chart[OF assms(1,2,2)] by (simp add: smooth-compat-commute)

```

```

lemma R-chart-in-atlas: R-chart x c ∈ atlas if x ∈ carrier c ∈ atlas
proof (rule maximal-atlas)
  show domain (R-chart x c) ⊆ carrier using R-action-in that by auto
  fix c' assume c' ∈ atlas
  with that(2) have ∞-smooth-compat c c' by (simp add: atlas-is-atlas)
  thus ∞-smooth-compat (R-chart x c) c'
  using smooth-compat-R-chart[OF that] by (simp add: ⟨c' ∈ atlas⟩)
qed

```

```

lemma right-action-automorphic: c-automorphism ∞ charts (R x) (R (invs x))
  if x ∈ carrier
proof (unfold-locales)
  fix y assume y ∈ carrier
  then obtain c1 where c1: c1 ∈ atlas y ∈ domain c1 using atlasE by blast
  let ?R = right-self-action x
  let ?Ri = right-self-action (invs x)

```

To find the second chart, for the codomain of $\lambda g'. \text{tms } g' (\text{invs } x)$, just shift the first chart across.

```

show ∃ c1 ∈ atlas. ∃ c2 ∈ atlas.
  y ∈ domain c1 ∧
  ?R ' domain c1 ⊆ domain c2 ∧
  smooth-on (codomain c1) (c2 ∘ ?R ∘ inv-chart c1)
proof (intro bexI conjI)
  let ?c2 = R-chart x c1

  show y ∈ domain c1 by (simp add: c1(2))

```

```

show  $c1 \in \text{atlas } ?c2 \in \text{atlas}$  by (simp add: R-chart-in-atlas c1(1) that)+
show  $(\lambda y. \text{tms } y \text{ (invs } x)) \text{ ' domain } c1 \subseteq \text{domain } ?c2$  by (simp add: c1(1) that)

have cong-to-id:  $(c1 \circ ?R_i \circ ?R \circ \text{inv-chart } c1) a = a$  if  $a \in \text{codomain } c1$  for  $a$ 
using R-action-invs(1)  $\langle x \in \text{carrier} \rangle c1(1)$  that by force
show smooth-on (codomain c1)  $(?c2 \circ ?R \circ \text{inv-chart } c1)$ 
using smooth-on-id smooth-on-cong cong-to-id
by (smt (verit, ccfv-threshold) R-chart-apply-chart c1(1) comp-apply open-codomain
that uminus-uminus)
qed

show  $\exists c1 \in \text{atlas}. \exists c2 \in \text{atlas}.$ 
 $y \in \text{domain } c1 \wedge$ 
 $?R_i \text{ ' domain } c1 \subseteq \text{domain } c2 \wedge$ 
 $\text{smooth-on (codomain } c1) (c2 \circ ?R_i \circ \text{inv-chart } c1)$ 
proof (intro bexI conjI)
let  $?c2 = \text{R-chart (invs } x) c1$ 

have [simp]:  $\text{invs } x \in \text{carrier}$  by (simp add: that uminus-mem)

show  $y \in \text{domain } c1$  by (simp add: c1(2))
show  $c1 \in \text{atlas } ?c2 \in \text{atlas}$  by (simp add: R-chart-in-atlas c1(1) that)+
show  $(\lambda g'. \text{tms } g' \text{ (invs (invs } x))) \text{ ' domain } c1 \subseteq \text{domain } ?c2$  by (simp add:
c1(1) that)

have 1:  $(c1 \circ ?R \circ ?R_i \circ \text{inv-chart } c1) a = a$ 
if  $a \in \text{codomain } c1$  for  $a$ 
using R-action-invs(2)  $\langle x \in \text{carrier} \rangle c1$  that by force
show smooth-on (codomain c1)  $(?c2 \circ ?R_i \circ \text{inv-chart } c1)$ 
apply (rule smooth-on-cong)
using 1 by (auto simp add: c1 uminus-uminus[OF that])
qed

{ fix  $y$  assume  $y \in \text{carrier}$ 
show  $\text{tms (tms } y \text{ (invs } x)) \text{ (invs (invs } x)) = y$ 
by (metis  $\langle y \in \text{carrier} \rangle$  add-assoc add-zeroR right-minus that uminus-mem)
show  $\text{tms (tms } y \text{ (invs (invs } x))) \text{ (invs } x) = y$ 
by (metis  $\langle y \in \text{carrier} \rangle$  add-assoc add-zeroR left-minus that uminus-mem) }
qed

lemma right-action-in-Diff: right-action  $x \in \text{Diff}$  if  $x \in \text{carrier}$ 
apply (intro DiffI automorphismI exI[where  $x = \text{right-self-action (invs } x)$ ])
subgoal using c-automorphism.c-automorphism-cong right-action-automorphic
that by fastforce
subgoal by (simp add: domIff order-class.order-eq-iff subset-iff)
done

end

```

9 Models/Instances

9.1 Euclidean Space

Euclidean spaces are dealt with at the start of the section “Differentiable Functions” in *Smooth-Manifolds.Differentiable-Manifold*. Therefore, this section is really just a “trivial” exercise to get used to things.

9.1.1 Euclidean Spaces are Lie groups under (+).

locale *euclidean-lie-group-add*
begin

abbreviation *C*
where $C \equiv \text{manifold-eucl.carrier}$

abbreviation *C-prod*
where $C\text{-prod} \equiv \text{manifold.carrier prod-charts-eucl}$

lemma *eucl-is-group*: *group-on-with* C (+) 0 (−) *uminus*

proof (*unfold group-on-with-def, intro conjI*)

show *monoid-on-with* C (+) 0

unfolding *monoid-on-with-def semigroup-add-on-with-def*

using *manifold-eucl-carrier*

by (*simp add: monoid-on-with-axioms.intro*)

show *group-on-with-axioms* C (+) 0 (−) *uminus*

unfolding *group-on-with-axioms-def*

using *manifold-eucl-carrier UNIV-I ab-group-add-class.ab-diff-conv-add-uminus add.left-inverse*

by *auto*

qed

lemma *prod-domain-codomain*: *domain prod-chart-eucl* = $C \times C$ *C* × *C* = *C-prod*
codomain prod-chart-eucl = $C \times C$

using *c-manifold-prod.domain-prod-chart [OF eucl-makes-product-manifold]*

apply *fastforce*

using *c-manifold-prod.prod-carrier eucl-makes-product-manifold*

apply *metis*

using *c-manifold-prod.codomain-prod-chart [OF eucl-makes-product-manifold]*

by *fastforce*

lemma *smooth-on-add-const*: *smooth-on* C ($\lambda a. a+b$)

proof –

have *sm-id*: *smooth-on* C ($\lambda a. a$)

by (*simp add: smooth-on-id*)

have *sm-add*: *smooth-on* C ($\lambda a. b$)

by (*simp add: smooth-on-const*)

show *smooth-on* C ($\lambda a. a+b$)

using *smooth-on-add [OF sm-id sm-add manifold.open-carrier]*

by simp
qed

lemma *smooth-binop-diff*:

fixes *tms*::'a \Rightarrow 'a \Rightarrow 'a::euclidean-space

assumes *smooth-on C-prod* ($\lambda(a,b). tms a b$)

shows *diff* ∞ *prod-charts-eucl charts-eucl* ($\lambda(x, y). tms x y$)

proof (*unfold diff-def diff-axioms-def, intro conjI allI impI*)

let *?prod* = *prod-charts-eucl*

let *?mult* = $\lambda(x, y). tms x y$

let *?c1* = *prod-chart-eucl*

let *?c2* = *chart-eucl*

let *?atl* = *manifold-eucl.atlas* ∞

let *?prod-atl* = *c-manifold.atlas prod-charts-eucl* ∞

fix *p*::'a \times 'a

assume *p* \in *manifold.carrier ?prod*

show $\exists c1 \in c\text{-manifold.atlas prod-charts-eucl } \infty. \exists c2 \in \text{manifold-eucl.atlas } \infty.$

$p \in \text{domain } c1 \wedge$

$(\lambda(x, y). tms x y) \text{ ' domain } c1 \subseteq \text{domain } c2 \wedge$

smooth-on (*codomain c1*) (*apply-chart c2* \circ ($\lambda(x, y). tms x y$) \circ *inv-chart c1*)

proof (*intro bexI, intro conjI*)

show *?c1* \in *?prod-atl*

by (rule *c-manifold.in-charts-in-atlas* [
OF *c-manifold-prod.c-manifold-atlas-product* [
OF *eucl-makes-product-manifold*
] *prod-chart-in-prod-charts*
])

show *?c2* \in *?atl*

using *c-manifold.in-charts-in-atlas* by simp

show *p* \in *domain ?c1*

by (*simp add: prod-domain-codomain*)

show $(\lambda(x, y). tms x y) \text{ ' domain } ?c1 \subseteq \text{domain } ?c2$

by simp

show *smooth-on* (*codomain ?c1*) (*apply-chart ?c2* \circ ($\lambda(x, y). tms x y$) \circ *inv-chart ?c1*)

using *map-fun-eucl-prod-id-f prod-domain-codomain assms*

by *metis*

qed

qed (*simp add: c-manifold-prod.c-manifold-atlas-product c-manifolds.intro eucl-makes-product-manifold manifold-eucl.c-manifold-axioms*)

lemma *smooth-unop-diff*:

fixes *invs*::'a \Rightarrow 'a::euclidean-space

assumes *smooth-on C invs*

shows *diff* ∞ *charts-eucl charts-eucl* *invs*

proof (*unfold diff-def diff-axioms-def, intro conjI allI impI*)

let *?c1* = *prod-chart-eucl*

let *?c2* = *chart-eucl*

let *?atl* = *manifold-eucl.atlas* ∞

```

fix x::'a
assume x ∈ manifold-eucl.carrier
show ∃ c1∈manifold-eucl.atlas ∞. ∃ c2∈manifold-eucl.atlas ∞.
  x ∈ domain c1 ∧
  invs ' domain c1 ⊆ domain c2 ∧
  smooth-on (codomain c1) (apply-chart c2 ∘ invs ∘ inv-chart c1)
proof (intro bexI conjI)
  show invs ' domain chart-eucl ⊆ domain ?c2
  by (simp add: image-subsetI)
  have manifold-eucl.carrier = codomain chart-eucl
  by simp
  thus smooth-on (codomain chart-eucl) (apply-chart chart-eucl ∘ invs ∘ inv-chart
chart-eucl)
  using assms map-fun-eucl-id-f
  by metis
  qed (simp+)
qed (simp add: manifold-eucl.self.c-manifolds-axioms)

```

```

lemma eucl-smooth-group-imp-lie-group:
  assumes is-group: group-on-with C tms tms-1 dvsn invs
  and smooth-mult: smooth-on C-prod (λ(a,b). tms a b)
  and smooth-inv: smooth-on C invs
  shows lie-group charts-eucl tms tms-1 dvsn invs
proof (unfold lie-group-def lie-group-axioms-def, (intro conjI))
  show c-manifold charts-eucl ∞
  using c-manifold-def by (simp add: c1-manifold-atlas-eucl)
  show group-on-with manifold-eucl.carrier tms tms-1 dvsn invs
  using is-group by simp
  show diff ∞ prod-charts-eucl charts-eucl (λ(a, b). tms a b)
  using smooth-binop-diff smooth-mult by auto
  show diff ∞ charts-eucl charts-eucl invs
  using smooth-unop-diff smooth-inv by simp
qed

```

Any Euclidean space is a Lie group under addition.

```

theorem lie-group-eucl: lie-group charts-eucl (+) 0 (-) uminus
by (rule eucl-smooth-group-imp-lie-group [OF eucl-is-group eucl-add-smooth eucl-um-smooth])

```

```

interpretation lie-group-eucl: lie-group charts-eucl (+) 0 (-) uminus
using lie-group-eucl .

```

end

9.2 The real numbers as a Lie group

```

lift-definition chart-real::(real, real) chart is
  (UNIV, UNIV, λx. x, λx. x)
by (auto simp: homeomorphism-def)

```

```

abbreviation charts-real ≡ {chart-real}

```

lemma *chart-real-is-eucl*: *charts-eucl = charts-real chart-eucl = chart-real*
by (*transfer, simp*)⁺

theorem *lie-group-real*: *lie-group charts-real (+) 0 (-) uminus*
using *euclidean-lie-group-add.lie-group-eucl chart-real-is-eucl* **by** *metis*

end

10 The Lie algebra of a Lie Group

theory *Lie-Algebra*

imports

Lie-Group

Manifold-Lie-Bracket

Smooth-Manifolds.Cotangent-Space

begin

sublocale *lie-group* \subseteq *smooth-manifold* **by** *unfold-locales*

locale *lie-algebra-morphism* =

src: *lie-algebra S1 scale1 bracket1* +

dest: *lie-algebra S2 scale2 bracket2* +

linear-on S1 S2 scale1 scale2 f

for *S1 S2*

and *scale1::'a::field* \Rightarrow *'b* \Rightarrow *'b::ab-group-add* **and** *scale2::'a::field* \Rightarrow *'c* \Rightarrow
'c::ab-group-add

and *bracket1* **and** *bracket2*

and *f* +

assumes *bracket-hom*: $\bigwedge X Y. X \in S1 \Longrightarrow Y \in S1 \Longrightarrow f (bracket1 X Y) =$
bracket2 (f X) (f Y)

Multiple isomorphic Lie algebras can be referred to as “the” Lie algebra \mathfrak{g} of a given Lie group G . One Lie algebra is already guaranteed to exist for any Lie group by virtue of *smooth-manifold ?charts* \Longrightarrow *lie-algebra (smooth-manifold.SVF ?charts) (*_R) lie-bracket-of-smooth-vector-fields*. We give an isomorphism between the subalgebra of *left-invariant* (smooth) vector fields and the tangent space at identity, and take the latter to be “the” Lie algebra \mathfrak{g} .

context *lie-group* **begin**

Some notation, for simplicity: the Lie group (or here, its carrier) is G , and the tangent space at the identity (the Lie algebra) is \mathfrak{g} .

notation *carrier* ($\langle G \rangle$)

definition *tangent-space-at-identity* ($\langle \mathfrak{g} \rangle$)

where *tangent-space-at-identity* = *tangent-space tms-one*

10.1 (Left-)invariant vector fields

A vector field X is invariant under some k -smooth map F if the vector assigned to a point $F(p)$ by X is the same as the vector assigned by (the push-forward under) F to the vector $X(p)$. Essentially, F and X “commute”.

definition (in c -manifold) *vector-field-invariant-under* :: 'a vector-field \Rightarrow ('a \Rightarrow 'a) \Rightarrow bool

(**infix** <invariant'-under> 80)

where X invariant-under $F \equiv \forall p \in \text{carrier}. \forall f \in \text{diff-fun-space}.$

$$X (F p) f = (\text{diff.push-forward } k \text{ charts } \text{charts } F) (X p) f$$

— TODO this could be in an instance of *diff* going from a manifold to itself, rather than *diffeomorphism*, i.e. an endomorphism rather than an automorphism.

definition (in c -automorphism) *invariant* :: 'a vector-field \Rightarrow bool

where $\text{invariant } X \equiv \forall p \in \text{carrier}. \forall g \in \text{src.diff-fun-space}. X (f p) g = \text{push-forward } (X p) g$

lemma (in c -automorphism) *invariant-simp*: $\text{src.vector-field-invariant-under } X f = \text{invariant } X$

unfolding $\text{src.vector-field-invariant-under-def}$ *invariant-def* **by** *simp*

lemma (in c -manifold) *vector-field-invariant-underD*: $X (F p) f = X p (\text{restrict0 carrier } (f \circ F))$

if X invariant-under F *diff* k charts charts F $p \in \text{carrier}$ $f \in \text{diff-fun-space}$

using *that* **by** (*auto simp: vector-field-invariant-under-def diff.push-forward-def*)

lemma (in c -manifold) *vector-field-invariant-underI*: X invariant-under F

if $\text{diff } k$ charts charts $F \wedge p f. p \in \text{carrier} \implies f \in \text{diff-fun-space} \implies X (F p) f = X p (\text{restrict0 carrier } (f \circ F))$

by (*simp add: vector-field-invariant-under-def diff.push-forward-def that*)

— Repeat notation from c -manifold $?charts ?k \implies c\text{-manifold.vector-field-invariant-under } ?charts ?k ?X ?F \equiv \forall p \in \text{manifold.carrier } ?charts. \forall f \in c\text{-manifold.diff-fun-space } ?charts ?k. ?X (?F p) f = \text{diff.push-forward } ?k ?charts ?charts ?F (?X p) f.$

notation *vector-field-invariant-under* (**infix** <invariant'-under> 80)

abbreviation $L\text{-invariant } X \equiv \forall p \in \text{carrier}. X \text{ invariant-under } (\mathcal{L} p)$

lemma $L\text{-invariantD}$ [*dest*]: $X (tms p q) f = X q (\text{restrict0 } G (f \circ (\mathcal{L} p)))$

if $L\text{-invariant } X$ $p \in G$ $q \in G$ $f \in \text{diff-fun-space}$

using $\text{vector-field-invariant-underD}$ *diff-tms(1)* **that** **by** *auto*

lemma $L\text{-invariantI}$ [*intro*]: $L\text{-invariant } X$

if $\wedge p q f. p \in \text{carrier} \implies q \in \text{carrier} \implies f \in \text{diff-fun-space} \implies X (tms p q) f = X q (\text{restrict0 carrier } (f \circ (\mathcal{L} p)))$

using *that* $\text{vector-field-invariant-underI}$ *diff-tms(1)* **by** *auto*

lemma *lie-bracket-left-invariant*:

assumes $L\text{-invariant } X$ *smooth-vector-field* X

and $L\text{-invariant } Y$ *smooth-vector-field* Y

```

shows L-invariant [X;Y] smooth-vector-field [X;Y]
proof
  fix p assume p: p ∈ G
  show vector-field-invariant-under [X;Y] ( $\mathcal{L}$  p)
  proof (intro vector-field-invariant-underI)
    fix q f
    assume q: q ∈ G and f: f ∈ diff-fun-space
    have 1: restrict0 G ((Z' f) ∘  $\mathcal{L}$  p) = Z'' (restrict0 G (f ∘  $\mathcal{L}$  p))
    if Z: L-invariant Z extensional0 carrier Z for Z
    proof
      fix t show restrict0 G ((Z'' f) ∘ tms p) t = Z t (restrict0 G (f ∘ tms p))
      apply (cases t ∈ G)
      subgoal
        using f p Z vector-field-invariant-underD[OF - - q smooth-vf-diff-fun-space]
        by (auto)
        using Z by (simp add: extensional0-outside)
      qed
      show [X;Y] (tms p q) f = [X;Y] q (restrict0 G (f ∘ tms p))
      unfolding lie-bracket-def
      using assms diff-tms(1) assms
      by (auto simp: 1 p f vector-field-invariant-underD[OF - - q smooth-vf-diff-fun-space]
smooth-vector-fieldE(2))
      qed (simp add: p diff-tms(1))
    qed (simp-all add: assms(2,4) lie-bracket-closed)

```

In fact, left-invariant smooth vector fields form a Lie subalgebra.

```

lemma subspace-of-left-invariant-svf:
  fixes  $\mathfrak{X}_{\mathcal{L}}$  defines  $\mathfrak{X}_{\mathcal{L}} \equiv \{X \in SVF. L\text{-invariant } X\}$ 
  shows subspace  $\mathfrak{X}_{\mathcal{L}}$ 
proof (unfold subspace-def, safe)
  interpret SVF: lie-algebra SVF scaleR lie-bracket-of-smooth-vector-fields
  using lie-algebra-of-smooth-vector-fields by simp

  have L-invariant 0
  apply (intro ballI vector-field-invariant-underI) by (simp-all add: diff-tms(1))
  thus 0 ∈  $\mathfrak{X}_{\mathcal{L}}$  unfolding assms(1) using SVF.m1.mem-zero by blast

  fix c and x
  assume x: x ∈  $\mathfrak{X}_{\mathcal{L}}$ 
  then have L-invariant (c *R x)
  apply (intro ballI vector-field-invariant-underI) using assms by (auto simp
add: diff-tms(1))
  thus c *R x ∈  $\mathfrak{X}_{\mathcal{L}}$  unfolding assms(1) using SVF.m1.mem-scale x assms by
blast

  fix y
  assume y: y ∈  $\mathfrak{X}_{\mathcal{L}}$ 
  then have L-invariant (x + y)
  apply (intro ballI vector-field-invariant-underI) using assms vector-field-invariant-underD

```

```

x by (auto simp: diff-tms(1))
  thus  $x + y \in \mathfrak{X}_{\mathcal{L}}$  unfolding assms(1) using SVF.m1.mem-add assms x y by
blast
qed

```

```

lemma lie-algebra-of-left-invariant-svf:
  fixes  $\mathfrak{X}_{\mathcal{L}}$  defines  $\mathfrak{X}_{\mathcal{L}} \equiv \{X. \text{smooth-vector-field } X \wedge L\text{-invariant } X\}$ 
  shows lie-algebra  $\mathfrak{X}_{\mathcal{L}}$  ( $*_R$ ) ( $\lambda X Y. [X; Y]$ )
proof -
  interpret SVF: lie-algebra SVF scaleR lie-bracket-of-smooth-vector-fields
  using lie-algebra-of-smooth-vector-fields by simp
  show ?thesis
  using assms subspace-of-left-invariant-svf by (auto intro: SVF.lie-subalgebra
simp: SVF.m1.implicit-subspace-with subspace-with lie-bracket-left-invariant
SVF-def)
qed

end

end

```

```

theory Classical-Groups

```

```

imports
  Lie-Group
  Linear-Algebra-More

```

```

begin

```

11 Matrix Groups

11.1 Entry Type

What would be a good type for the entries of our matrices? Ideally, I would be able to talk about matrices over reals \mathbb{R} , the complex numbers \mathbb{C} , and the quaternionic skew-field \mathbb{H} . This is hard: only algebras and inner product spaces over \mathbb{R} are well-supported in Isabelle's Main.

For now, for simplicity, I will work with real matrices only. Alternatively, one could try to characterise the type class containing \mathbb{R} , \mathbb{C} , and \mathbb{H} only. Below is a first attempt to maintain at least some generality. I give some trivial type instantiations, as a basic check.

However, locales are the way to go, in my opinion.

```

class real-normed-eucl = real-normed-field + euclidean-space

```

instance *real-normed-eucl* \subseteq *euclidean-space* **by** *standard*
instance *real-normed-eucl* \subseteq *real-normed-field* **by** *standard*
instance *real-normed-eucl* \subseteq *topological-space* **by** *standard*

instance *real-normed-eucl* \subseteq *comm-ring* **by** *standard*
instance *real-normed-eucl* \subseteq *comm-ring-1* **by** *standard*
instance *real-normed-eucl* \subseteq *real-algebra-1* **by** *standard*

instance *vec* :: (*real-normed-eucl*, *finite*) *topological-space* **by** *standard*
instance *vec* :: (*real-normed-eucl*, *finite*) *euclidean-space* **by** *standard*

instance *real* :: *real-normed-eucl* **by** *standard*
instance *complex* :: *real-normed-eucl* **by** *standard*

11.2 Mat(**n**, **F**)

The set of all '*n*-vectors over a *topological-space* is a *topological-space*: this is proved in *Finite-Cartesian-Product*. Similar for vectors over a *euclidean-space*. Therefore, a vector of vectors over a topological space (i.e. a matrix) is also a topological space. We can thus define the identity as a chart; this is not superbly useful, but serves as a template for charts for the multiplicative matrix groups later on.

lift-definition *chart-mat*::(*'a*::*real-normed-eucl*,*'n*::*finite*)*square-matrix*, (*'a*,*'n*)*square-matrix*)*chart*
is (*UNIV*, *UNIV*, $\lambda m. m$, $\lambda m. m$)
by (*auto simp: homeomorphism-def*)

11.3 GL(**n**, **F**)

We define polymorphic abbreviations for the carrier set of the general linear group as a matrix group over a commutative ring. This group can be considered as the automorphism group on arbitrary modules of non-commutative rings too, but one loses the isomorphism with matrices, and I'm mostly interested in much more specific general linear groups anyway (namely, over real and complex numbers). Using commutative rings (with 1) also means that determinants play nicely.

abbreviation *in-GL*::(*'a*::*comm-ring-1*,*'n*::*finite*)*square-matrix* \Rightarrow *bool*
where *in-GL* \equiv *invertible*
abbreviation *GL* **where** *GL* \equiv *Collect in-GL*

As an example for making the polymorphic *GL* concrete, we specify the general linear group in four real/complex dimensions.

abbreviation *GL_{R4}*::(*real*,4)*square-matrix set* **where** *GL_{R4}* \equiv *GL*
abbreviation *GL_{C4}*::(*complex*,4)*square-matrix set* **where** *GL_{C4}* \equiv *GL*

PROBLEM: the inner product on the LHS is real, not complex, which is why the commented line (involving complex multiplication) cannot work (it only passes type checking because *complex-of-real* is a coercion).

lemma

assumes $x \in GL_{C4}$

shows $((row\ i\ x \cdot row\ i\ x)::real) = (\sum_{j \in UNIV}. (row\ i\ x)\$j \cdot (row\ i\ x)\$j)$
by (*simp add: inner-vec-def*)

We now define the chart that makes $GL(n,F)$ a Lie group. Since a chart is a homeomorphism, we first need to show that GL is an open set. Notice this GL is already restricted to have much more powerful entries, since we require topology (continuity) now.

lemma *GL-preimage-det*: $det - ' (UNIV - \{0::'a::real-normed-eucl\}) = GL$

proof (*safe*)

fix $x::('a::real-normed-eucl, 'n::finite)\ square-matrix$

assume *in-GL x*

then show $x \in det - ' (UNIV - \{0\})$

using *invertible-det-nz* **by** *auto*

next

fix $x::('a::real-normed-eucl, 'n::finite)\ square-matrix$

assume $det\ x \neq 0$

then show *in-GL x*

by (*simp add: invertible-det-nz*)

qed

lemma *open-GL*: $open\ (GL::('a::real-normed-eucl, 'n::finite)\ square-matrix\ set)$

using *open-vimage continuous-on-det GL-preimage-det*

by (*metis open-UNIV open-delete*)

lift-definition *chart-GL*:: $(('a::real-normed-eucl, 'n::finite)\ square-matrix, ('a, 'n)\ square-matrix)\ chart$

is $(GL, GL, \lambda m. m, \lambda m. m)$

by (*auto simp: homeomorphism-def open-GL*)

lift-definition *real-chart-GL*:: $((real, 'n::finite)\ square-matrix, (real, 'n)\ square-matrix)\ chart$

is $(GL, GL, \lambda m. m, \lambda m. m)$

by (*auto simp: homeomorphism-def open-GL*)

lemma *transfer-GL* [*simp*]:

shows $domain\ chart-GL = GL$

and $codomain\ chart-GL = GL$

and $apply-chart\ chart-GL = (\lambda x. x)$

and $inv-chart\ chart-GL = (\lambda x. x)$

by (*transfer, simp*)⁺

abbreviation *charts-GL* **where** $charts-GL \equiv \{chart-GL\}$

abbreviation *real-charts-GL* **where** $real-charts-GL \equiv \{real-chart-GL\}$

interpretation *manifold-GL*: *c-manifold charts-GL k*
using *smooth-compat-refl* **by** (*unfold-locales*, *simp*)

abbreviation *prod-chart-GL* :: (*'a::real-normed-eucl*, *'b::finite*)*square-matrix* ×
(*'a*, *'b*)*square-matrix*, (*'a*, *'b*)*square-matrix* × (*'a*, *'b*)*square-matrix*) *chart*
where *prod-chart-GL* ≡ *c-manifold-prod.prod-chart chart-GL chart-GL*
abbreviation *prod-charts-GL* :: (*'a::real-normed-eucl*, *'b::finite*)*square-matrix* ×
(*'a*, *'b*)*square-matrix*, (*'a*, *'b*)*square-matrix* × (*'a*, *'b*)*square-matrix*) *chart set*
where *prod-charts-GL* ≡ *c-manifold-prod.prod-charts charts-GL charts-GL*

interpretation *prod-manifold-GL*: *c-manifold-prod k*
charts-GL::('a::real-normed-eucl,'n::finite)square-matrix, ('a,'n)square-matrix) chart
set
charts-GL::('a::real-normed-eucl,'n::finite)square-matrix, ('a,'n)square-matrix) chart
set
unfolding *c-manifold-prod-def* **apply** (*simp add: manifold-GL.c-manifold-axioms*)
done

abbreviation *prod-GL-carrier* ≡ *manifold.carrier prod-manifold-GL.prod-charts*
abbreviation *prod-GL-atlas* ≡ *c-manifold.atlas prod-manifold-GL.prod-charts* ∅

lemma *transfer-prod-GL [simp]*:
shows *domain prod-chart-GL = GL × GL*
and *codomain prod-chart-GL = GL × GL*
and *apply-chart prod-chart-GL = (λx. x)*
and *inv-chart prod-chart-GL = (λx. x)*
using *c-manifold-prod.domain-prod-chart c-manifold-prod.codomain-prod-chart*
c-manifold-prod.apply-prod-chart c-manifold-prod.inv-chart-prod-chart trans-
fer-GL
by *auto*

lemma *manifold-GL-carrier [simp]*: *manifold-GL.carrier = GL*
by (*simp add: manifold-GL.carrier-def*)

lemma *prod-manifold-GL-carrier [simp]*: *prod-GL-carrier = GL × GL*
using *prod-manifold-GL.prod-carrier* **by** *auto*

The following lemma basically just does unfolding and type checking.
Possibly useful once general results for *charts-GL* need to be specified down
to *real-charts-GL*.

lemma *real-GL-is-a-GL*:
shows *real-chart-GL = chart-GL*
and *real-charts-GL = charts-GL*
and *manifold.carrier (c-manifold-prod.prod-charts real-charts-GL real-charts-GL)*
= *prod-GL-carrier*
unfolding *chart-GL-def real-chart-GL-def* **by** *simp+*

```

lemma mult-closed-on-GL:
  fixes f-mult :: ('a,'b)square-matrix × ('a,'b)square-matrix
    ⇒ ('a::comm-ring-1, 'b::finite) square-matrix
  defines f-mult: f-mult ≡ (λ(x, y). x ** y)
  shows f-mult ' (GL × GL) ⊆ GL
proof
  fix x
  assume x ∈ f-mult ' (GL × GL)
  then obtain y z::('a,'b)square-matrix where x = y**z invertible y invertible z
    using f-mult by auto
  then show x ∈ GL
    by (simp add: invertible-mult)
qed

lemma GL-group-mult-right-div:
  shows group-on-with (domain chart-GL) (**) (mat 1) (λm1 m2. m1 ** matrix-inv
m2) matrix-inv
  apply unfold-locales
  apply (simp-all add: matrix-mul-assoc invertible-mult invertible-mat-1 invertible-matrix-inv)
  by (simp add: matrix-inv-def invertible-right-inverse matrix-left-right-inverse
verit-sko-ex-indirect)

lemma smooth-on-proj: smooth-on prod-GL-carrier fst smooth-on prod-GL-carrier
snd
  using smooth-on-fst [OF smooth-on-id manifold.open-carrier] apply blast
  using smooth-on-snd [OF smooth-on-id manifold.open-carrier] by blast

lemma mult-smooth-on-real-GL:
  fixes f-mult :: (real,'n)square-matrix × (real,'n)square-matrix ⇒ (real,'n::finite)square-matrix
  defines f-mult: f-mult ≡ (λ(x, y). x ** y)
  shows smooth-on (GL × GL) f-mult
proof (unfold f-mult, simp add: case-prod-beta', intro smooth-on-matrix-mult)
  — Isabelle doesn't seem to infer types for GL and prod-GL-carrier below, even
  though they should be clear from being accepted in “show” statements (i.e. they
  should be inferred from having to match the types in the lemma's goal).
  let ?GL = GL::(real,'n)square-matrix set
  show smooth-on (?GL × ?GL) fst
    using smooth-on-proj(1) by simp
  show smooth-on (?GL × ?GL) snd
    using smooth-on-proj(2) by simp
  show open (?GL × ?GL)
    using manifold.open-carrier[of prod-charts-GL] prod-manifold-GL-carrier by
simp
qed

```

lemma *mult-smooth-on-GL-expanded*:
assumes $x \in \text{prod-GL-carrier}$
shows $x \in \text{domain prod-chart-GL}$
and $(\lambda(x, y). x ** y) \text{ ' domain prod-chart-GL } \subseteq \text{ domain chart-GL}$
and $\text{smooth-on (codomain prod-chart-GL) (apply-chart chart-GL } \circ (\lambda(x, y). x$
 $** y) \circ \text{inv-chart prod-chart-GL)}$
using *assms apply fastforce*
apply (*simp add: mult-closed-on-GL*)
apply (*simp add: fun.map-ident*)
using *mult-smooth-on-real-GL* — only for real entries
oops

lemma *mult-smooth-on-real-GL-expanded*:
fixes $f\text{-mult} :: (\text{real}, 'n)\text{square-matrix} \times (\text{real}, 'n)\text{square-matrix} \Rightarrow (\text{real}, 'n::\text{finite})\text{square-matrix}$
and $x :: (\text{real}, 'n)\text{square-matrix} \times (\text{real}, 'n)\text{square-matrix}$
defines $f\text{-mult}: f\text{-mult} \equiv (\lambda(x, y). x ** y)$
assumes $x \in \text{prod-GL-carrier}$
shows $x \in \text{domain prod-chart-GL}$
and $f\text{-mult} \text{ ' domain prod-chart-GL } \subseteq \text{ domain chart-GL}$
and $\text{smooth-on (codomain prod-chart-GL) (apply-chart chart-GL } \circ f\text{-mult } \circ$
 $\text{inv-chart prod-chart-GL)}$
proof –
show $x \in \text{domain prod-chart-GL}$
using *assms by fastforce*
show $f\text{-mult} \text{ ' domain prod-chart-GL } \subseteq \text{ domain chart-GL}$
by (*simp add: f-mult mult-closed-on-GL*)
show $\text{smooth-on (codomain prod-chart-GL) (apply-chart chart-GL } \circ f\text{-mult } \circ$
 $\text{inv-chart prod-chart-GL)}$
apply (*simp add: fun.map-ident*)
by (*simp add: f-mult mult-smooth-on-real-GL*)
qed

theorem *real-GL-Lie-group*: $\text{lie-group real-charts-GL (**) (mat 1) } (\lambda m_1 m_2. m_1$
 $** (\text{matrix-inv } m_2)) \text{ matrix-inv}$
proof (*intro group-manifold-imp-lie-group2*)
let $?div = \lambda m_1 m_2. m_1 ** \text{matrix-inv } m_2$
let $?prod = c\text{-manifold-prod.prod-charts real-charts-GL real-charts-GL}$
show $c\text{-manifold real-charts-GL } \infty$
by (*simp add: manifold-GL.c-manifold-axioms real-GL-is-a-GL(2)*)
show $\text{group-on-with } (\bigcup (\text{domain ' real-charts-GL})) (**) (\text{mat } 1) ?div \text{matrix-inv}$
using *GL-group-mult-right-div real-GL-is-a-GL(2)*
by (*metis (mono-tags, lifting) manifold.carrier-def manifold-GL-carrier trans-*
 $\text{fer-GL}(1)$)
show $\text{diff-axioms } \infty ?prod \text{real-charts-GL } (\lambda(a, b). a ** b)$
proof (*unfold-diff-axioms; unfold real-GL-is-a-GL(1,2) prod-manifold-GL-carrier*)
fix $x :: ((\text{real}, 'b) \text{vec}, 'b) \text{vec} \times ((\text{real}, 'b) \text{vec}, 'b) \text{vec}$
assume $x\text{-in}: x \in \text{GL} \times \text{GL}$


```

show  $x \in \text{domain prod-chart-GL}$ 
  using  $x\text{-in}$  by simp
show  $\text{real-chart-GL} \in \text{manifold-GL.atlas} \infty$ 
  by (simp add: manifold-GL.in-charts-in-atlas real-GL-is-a-GL(1))
show  $\text{prod-chart-GL} \in \text{prod-GL-atlas}$ 
  by (simp add: prod-manifold-GL.prod-chart-in-atlas)
show  $\text{mult-maps-domain: } (\lambda(x, y). x ** y) \text{ ' domain prod-chart-GL} \subseteq \text{domain}$ 
 $\text{real-chart-GL}$ 
  using  $x\text{-in mult-smooth-on-real-GL-expanded(2)}$ 
  by (simp add: mult-closed-on-GL real-GL-is-a-GL(1))
show  $\text{smooth-on (codomain prod-chart-GL) (}$ 
   $\text{apply-chart real-chart-GL} \circ (\lambda(x, y). x ** y) \circ \text{inv-chart prod-chart-GL)}$ 
  using  $\text{mult-smooth-on-real-GL-expanded(3) } x\text{-in real-GL-is-a-GL(1)}$ 
  by (metis prod-manifold-GL-carrier smooth-on-open-subsetsI transfer-prod-GL(2))
qed
show  $\text{diff-axioms} \infty \text{ real-charts-GL real-charts-GL matrix-inv}$ 
proof (unfold-diff-axioms; unfold real-GL-is-a-GL(1,2) manifold-GL-carrier)
  fix  $x :: (\text{real, 'b}) \text{vec, 'b}) \text{vec}$ 
  assume  $x\text{-in: } x \in \text{GL}$ 

```

— Cheeky: "cast up" the real matrix x to be in the domain of *chart-GL*, rather than *real-chart-GL*. This makes proofs easier for sledgehammer wherever they involve lemmas about GL in general.

```

  show  $x \in \text{domain real-chart-GL}$ 
    using  $x\text{-in}$  by (simp add: real-GL-is-a-GL(1))
  show  $\text{chart-GL} \in \text{manifold-GL.atlas} \infty$ 
    by (simp add: manifold-GL.in-charts-in-atlas)
  show  $\text{real-chart-GL} \in \text{manifold-GL.atlas} \infty$ 
    by (simp add: manifold-GL.in-charts-in-atlas real-GL-is-a-GL(1))
  show  $\text{mult-maps-domain: matrix-inv ' domain real-chart-GL} \subseteq \text{domain chart-GL}$ 
    by (simp add: image-subset-iff invertible-matrix-inv real-GL-is-a-GL(1))
  have  $1: (\lambda x. x) \circ \text{matrix-inv} \circ (\lambda x. x) = \text{matrix-inv}$ 
    by auto
  show  $\text{smooth-on (codomain real-chart-GL) (apply-chart chart-GL} \circ \text{matrix-inv}$ 
 $\circ \text{inv-chart real-chart-GL)}$ 
    using  $\text{smooth-on-matrix-inv[OF - open-GL]}$  by (simp add: real-GL-is-a-GL(1))
1)
qed
qed

```

```

corollary  $\text{real-GL-Lie-grp: lie-grp real-charts-GL (**) (mat 1)}$ 
  using  $\text{lie-group-imp-lie-grp[OF real-GL-Lie-group]}$  .

```

end

References

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