

# The Laws of Large Numbers

Manuel Eberl

## Abstract

The Law of Large Numbers states that, informally, if one performs a random experiment  $X$  many times and takes the average of the results, that average will be very close to the expected value  $E[X]$ .

More formally, let  $(X_i)_{i \in \mathbb{N}}$  be a sequence of independently identically distributed random variables whose expected value  $E[X_1]$  exists. Denote the running average of  $X_1, \dots, X_n$  for  $\bar{X}_n$ . Then:

- The Weak Law of Large Numbers states that  $\bar{X}_n \rightarrow E[X_1]$  in probability for  $n \rightarrow \infty$ , i.e.  $\mathcal{P}(|\bar{X}_n - E[X_1]| > \varepsilon) \rightarrow 0$  for  $n \rightarrow \infty$  for any  $\varepsilon > 0$ .
- The Strong Law of Large Numbers states that  $\bar{X}_n \rightarrow E[X_1]$  almost surely for  $n \rightarrow \infty$ , i.e.  $\mathcal{P}(\bar{X}_n \rightarrow E[X_1]) = 1$ .

In this entry, I formally prove the strong law and from it the weak law. The approach used for the proof of the strong law is a particularly quick and slick one based on ergodic theory, which was formalised by Gouëzel in another AFP entry.

## Contents

<b>1</b>	<b>The Laws of Large Numbers</b>	<b>1</b>
1.1	The strong law . . . . .	2
1.2	The weak law . . . . .	3
1.3	Example . . . . .	4

## 1 The Laws of Large Numbers

```
theory Laws-of-Large-Numbers
imports Ergodic-Theory.Shift-Operator
begin
```

We prove the strong law of large numbers in the following form: Let  $(X_i)_{i \in \mathbb{N}}$  be a sequence of i.i.d. random variables over a probability space  $M$ . Further assume that the expected value  $E[X_0]$  of  $X_0$  exists. Then the sequence of random variables

$$\bar{X}_n = \frac{1}{n} \sum_{i=0}^n X_i$$

of running averages almost surely converges to  $E[X_0]$ . This means that

$$\mathcal{P}[\bar{X}_n \longrightarrow E[X_0]] = 1 .$$

We start with the strong law.

## 1.1 The strong law

The proof uses Birkhoff's Theorem from Gouëzel's formalisation of ergodic theory [1] and the fact that the shift operator  $T(x_1, x_2, x_3, \dots) = (x_2, x_3, \dots)$  is ergodic. This proof can be found in various textbooks on probability theory/ergodic theory, e.g. the ones by Krengel [2, p. 24] and Simmonet [3, Chapter 15, pp. 311–325].

**theorem** (in *prob-space*) *strong-law-of-large-numbers-iid*:

**fixes**  $X :: \text{nat} \Rightarrow 'a \Rightarrow \text{real}$

**assumes** *indep*: *indep-vars*  $(\lambda \cdot \text{borel}) X \text{ UNIV}$

**assumes** *distr*:  $\bigwedge i. \text{distr } M \text{ borel } (X i) = \text{distr } M \text{ borel } (X 0)$

**assumes** *L1*: *integrable*  $M (X 0)$

**shows** *AE*  $x$  in  $M. (\lambda n. (\sum_{i < n} X i x) / n) \longrightarrow \text{expectation } (X 0)$

**proof** –

We adopt a more explicit view of  $M$  as a countably infinite product of i.i.d. random variables, indexed by the natural numbers:

**define**  $M' :: (\text{nat} \Rightarrow \text{real}) \text{ measure}$  **where**  $M' = \text{Pi}_M \text{ UNIV } (\lambda i. \text{distr } M \text{ borel } (X i))$

**have** [*measurable*]: *random-variable*  $\text{borel } (X i)$  **for**  $i$

**using** *indep* **by** (*auto simp: indep-vars-def*)

**have**  $M'\text{-eq}$ :  $M' = \text{distr } M (\text{Pi}_M \text{ UNIV } (\lambda i. \text{borel})) (\lambda x. \lambda i \in \text{UNIV}. X i x)$

**using** *indep unfolding*  $M'\text{-def}$  **by** (*subst (asm) indep-vars-iff-distr-eq-PiM*)

*auto*

**have** *space- $M'$* : *space*  $M' = \text{UNIV}$

**by** (*simp add:  $M'\text{-def}$  space- $\text{Pi}_M$* )

**have** *sets- $M'$*  [*measurable-cong*]: *sets*  $M' = \text{sets } (\text{Pi}_M \text{ UNIV } (\lambda i. \text{borel}))$

**by** (*simp add:  $M'\text{-eq}$* )

**interpret**  $M'$ : *prob-space*  $M'$

**unfolding**  $M'\text{-eq}$  **by** (*intro prob-space-distr*) *auto*

We introduce a shift operator that forgets the first variable in the sequence.

**define**  $T :: (\text{nat} \Rightarrow \text{real}) \Rightarrow (\text{nat} \Rightarrow \text{real})$  **where**

$T = (\lambda f. f \circ \text{Suc})$

**have** *funpow- $T$* :  $(T \text{ ^^ } i) = (\lambda f. f \circ (\lambda n. n + i))$  **for**  $i$

**by** (*induction i*) (*auto simp:  $T\text{-def}$* )

**interpret**  $T$ : *shift-operator-ergodic*  $\text{distr } M \text{ borel } (X 0) T M'$

**proof** –

**interpret**  $X0$ : *prob-space*  $\text{distr } M \text{ borel } (X 0)$

**by** (*rule prob-space-distr*) *auto*

```

show shift-operator-ergodic (distr M borel (X 0))
  by unfold-locales
show  $M' \equiv Pi_M UNIV (\lambda-. distr M borel (X 0))$ 
  unfolding M'-def by (subst distr)
qed (simp-all add: T-def)

have [intro]: integrable M' ( $\lambda f. f 0$ )
  unfolding M'-eq by (subst integrable-distr-eq) (use L1 in auto)
have AE f in M'. ( $\lambda n. T.birkhoff-sum (\lambda f. f 0) n f / real n$ )
   $\longrightarrow$  real-cond-exp M' T.Invariants ( $\lambda f. f 0$ ) f
  by (rule T.birkhoff-theorem-AE-nonergodic) auto
moreover have AE x in M'. real-cond-exp M' T.Invariants ( $\lambda f. f 0$ ) x =
  M'.expectation ( $\lambda f. f 0$ ) / M'.prob (space M')
  by (intro T.Invariants-cond-exp-is-integral-fmpt) auto
ultimately have AE f in M'. ( $\lambda n. T.birkhoff-sum (\lambda f. f 0) n f / real n$ )
   $\longrightarrow$  M'.expectation ( $\lambda f. f 0$ )
  by eventually-elim (simp-all add: M'.prob-space)
also have M'.expectation ( $\lambda f. f 0$ ) = expectation (X 0)
  unfolding M'-eq by (subst integral-distr) simp-all
also have T.birkhoff-sum ( $\lambda f. f 0$ ) = ( $\lambda n f. sum f \{..<n\}$ )
  by (intro ext) (simp-all add: T.birkhoff-sum-def funpow-T)
finally show ?thesis
  unfolding M'-eq by (subst (asm) AE-distr-iff) simp-all
qed

```

## 1.2 The weak law

To go from the strong law to the weak one, we need the fact that almost sure convergence implies convergence in probability. We prove this for sequences of random variables here.

**lemma** (*in prob-space*) *AE-convergence-imp-convergence-in-prob*:

**assumes** [*measurable*]:  $\bigwedge i. random-variable borel (X i) random-variable borel Y$

**assumes** *AE*: *AE x in M. ( $\lambda i. X i x$ )  $\longrightarrow$  Y x*

**assumes**  $\varepsilon > (0 :: real)$

**shows**  $(\lambda i. prob \{x \in space M. |X i x - Y x| > \varepsilon\}) \longrightarrow 0$

**proof** –

**define** *A* **where**  $A = (\lambda i. \{x \in space M. |X i x - Y x| > \varepsilon\})$

**define** *B* **where**  $B = (\lambda n. (\bigcup i \in \{n.. \}. A i))$

**have** [*measurable*]: *A i*  $\in$  *sets M* *B i*  $\in$  *sets M* **for** *i*

**unfolding** *A-def B-def* **by** *measurable*

**have** *AE x in M. x*  $\notin$   $(\bigcap i. B i)$

**using** *AE* **unfolding** *B-def A-def*

**by** *eventually-elim*

(*use  $\langle \varepsilon > 0 \rangle$  in  $\langle fastforce simp: tendsto-iff dist-norm eventually-at-top-linorder \rangle$ )*

**hence**  $(\bigcap i. B i) \in null-sets M$

**by** (*subst AE-iff-null-sets*) *auto*

**show**  $(\lambda i. prob (A i)) \longrightarrow 0$

```

proof (rule Lim-null-comparison)
  have  $(\lambda i. \text{prob } (B \ i)) \longrightarrow \text{prob } (\bigcap i. B \ i)$ 
  proof (rule finite-Lim-measure-decseq)
    show decseq  $B$ 
    by (rule decseq-SucI) (force simp: B-def)
  qed auto
  also have  $\text{prob } (\bigcap i. B \ i) = 0$ 
    using  $\langle (\bigcap i. B \ i) \in \text{null-sets } M \rangle$  by (simp add: measure-eq-0-null-sets)
  finally show  $(\lambda i. \text{prob } (B \ i)) \longrightarrow 0$  .
next
  have  $\text{prob } (A \ n) \leq \text{prob } (B \ n)$  for  $n$ 
    unfolding B-def by (intro finite-measure-mono) auto
  thus  $\forall_F n \text{ in at-top. norm } (\text{prob } (A \ n)) \leq \text{prob } (B \ n)$ 
    by (intro always-eventually) auto
  qed
qed

```

The weak law is now a simple corollary: we again have the same setting as before. The weak law now states that  $\bar{X}_n$  converges to  $E[X_0]$  in probability. This means that for any  $\varepsilon > 0$ , the probability that  $|\bar{X}_n - X_0| > \varepsilon$  vanishes as  $n \rightarrow \infty$ .

```

corollary (in prob-space) weak-law-of-large-numbers-iid:
  fixes  $X :: \text{nat} \Rightarrow 'a \Rightarrow \text{real}$  and  $\varepsilon :: \text{real}$ 
  assumes indep: indep-vars  $(\lambda \cdot. \text{borel } X \ UNIV$ 
  assumes distr:  $\bigwedge i. \text{distr } M \ \text{borel } (X \ i) = \text{distr } M \ \text{borel } (X \ 0)$ 
  assumes L1: integrable  $M \ (X \ 0)$ 
  assumes  $\varepsilon > 0$ 
  shows  $(\lambda n. \text{prob } \{x \in \text{space } M. |(\sum i < n. X \ i \ x) / n - \text{expectation } (X \ 0)| > \varepsilon\})$ 
 $\longrightarrow 0$ 
proof (rule AE-convergence-imp-convergence-in-prob)
  show AE  $x \ \text{in } M. (\lambda n. (\sum i < n. X \ i \ x) / n) \longrightarrow \text{expectation } (X \ 0)$ 
    by (rule strong-law-of-large-numbers-iid) fact+
next
  have [measurable]: random-variable borel  $(X \ i)$  for  $i$ 
    using indep by (auto simp: indep-vars-def)
  show random-variable borel  $(\lambda x. (\sum i < n. X \ i \ x) / \text{real } n)$  for  $n$ 
    by measurable
qed (use  $\langle \varepsilon > 0 \rangle$  in simp-all)

```

**end**

### 1.3 Example

```

theory Laws-of-Large-Numbers-Example
  imports Laws-of-Large-Numbers
begin

```

As an example, we apply the strong law to the proportion of successes in an independent sequence of coin flips with success probability  $p$ . We will show

that proportion of successful coin flips among the first  $n$  attempts almost surely converges to  $p$  as  $n \rightarrow \infty$ .

**lemma** (in *prob-space*) *indep-vars-iff-distr-eq-PiM'*:

**fixes**  $I :: 'i \text{ set}$  and  $X :: 'a \Rightarrow 'b$

**assumes**  $I \neq \{\}$

**assumes**  $rv: \bigwedge i. i \in I \implies \text{random-variable } (M' i) (X i)$

**shows**  $\text{indep-vars } M' X I \longleftrightarrow$

$$\text{distr } M (\Pi_M i \in I. M' i) (\lambda x. \lambda i \in I. X i x) = (\Pi_M i \in I. \text{distr } M (M' i)$$

$(X i)$ )

**proof** –

**from** *assms* **obtain**  $j$  **where**  $j: j \in I$

**by** *auto*

**define**  $N'$  **where**  $N' = (\lambda i. \text{if } i \in I \text{ then } M' i \text{ else } M' j)$

**define**  $Y$  **where**  $Y = (\lambda i. \text{if } i \in I \text{ then } X i \text{ else } X j)$

**have**  $rv: \text{random-variable } (N' i) (Y i)$  **for**  $i$

**using**  $j$  **by** (*auto simp: N'-def Y-def intro: assms*)

**have**  $\text{indep-vars } M' X I = \text{indep-vars } N' Y I$

**by** (*intro indep-vars-cong*) (*auto simp: N'-def Y-def*)

**also have**  $\dots \longleftrightarrow \text{distr } M (\Pi_M i \in I. N' i) (\lambda x. \lambda i \in I. Y i x) = (\Pi_M i \in I. \text{distr } M (N' i) (Y i))$

**by** (*intro indep-vars-iff-distr-eq-PiM rv assms*)

**also have**  $(\Pi_M i \in I. N' i) = (\Pi_M i \in I. M' i)$

**by** (*intro PiM-cong*) (*simp-all add: N'-def*)

**also have**  $(\lambda x. \lambda i \in I. Y i x) = (\lambda x. \lambda i \in I. X i x)$

**by** (*simp-all add: Y-def fun-eq-iff*)

**also have**  $(\Pi_M i \in I. \text{distr } M (N' i) (Y i)) = (\Pi_M i \in I. \text{distr } M (M' i) (X i))$

**by** (*intro PiM-cong distr-cong*) (*simp-all add: N'-def Y-def*)

**finally show** *?thesis* .

**qed**

**lemma** *indep-vars-PiM-components*:

**assumes**  $\bigwedge i. i \in A \implies \text{prob-space } (M i)$

**shows**  $\text{prob-space.indep-vars } (PiM A M) M (\lambda i f. f i) A$

**proof** (*cases*  $A = \{\}$ )

**case** *False*

**have**  $\text{distr } (PiM A M) (PiM A M) (\lambda x. \text{restrict } x A) = \text{distr } (PiM A M) (PiM A M) (\lambda x. x)$

**by** (*intro distr-cong*) (*auto simp: restrict-def space-PiM PiE-def extensional-def Pi-def*)

**also have**  $\dots = PiM A M$

**by** *simp*

**also have**  $\dots = PiM A (\lambda i. \text{distr } (PiM A M) (M i) (\lambda f. f i))$

**by** (*intro PiM-cong refl, subst distr-PiM-component*) (*auto simp: assms*)

**finally show** *?thesis*

**by** (*subst prob-space.indep-vars-iff-distr-eq-PiM'*) (*simp-all add: prob-space-PiM assms False*)

**next**

```

case True
interpret prob-space PiM A M
  by (intro prob-space-PiM assms)
show ?thesis
  unfolding indep-vars-def indep-sets-def by (auto simp: True)
qed

```

```

lemma indep-vars-PiM-components':
  assumes  $\bigwedge i. i \in A \implies \text{prob-space } (M \ i)$ 
  assumes  $\bigwedge i. i \in A \implies g \ i \in M \ i \rightarrow_M N \ i$ 
  shows prob-space.indep-vars (PiM A M) N ( $\lambda i f. g \ i \ (f \ i)$ ) A
  by (rule prob-space.indep-vars-compose2[OF prob-space-PiM indep-vars-PiM-components])
    (use assms in simp-all)

```

```

lemma integrable-bernoulli-pmf [intro]:
  fixes f :: bool  $\Rightarrow$  'a :: {banach, second-countable-topology}
  shows integrable (bernoulli-pmf p) f
  by (rule integrable-measure-pmf-finite) auto

```

```

lemma expectation-bernoulli-pmf:
  fixes f :: bool  $\Rightarrow$  'a :: {banach, second-countable-topology}
  assumes p: p  $\in$  {0..1}
  shows measure-pmf.expectation (bernoulli-pmf p) f = p *R f True + (1 - p)
  *R f False
  using p by (subst integral-measure-pmf[of UNIV]) (auto simp: UNIV-bool)

```

```

experiment
  fixes p :: real
  assumes p: p  $\in$  {0..1}
begin

```

```

definition M :: (nat  $\Rightarrow$  bool) measure
  where M = ( $\Pi_M \ i \in (UNIV :: \text{nat set}). \text{measure-pmf } (\text{bernoulli-pmf } p)$ )

```

```

definition X :: nat  $\Rightarrow$  (nat  $\Rightarrow$  bool)  $\Rightarrow$  real
  where X = ( $\lambda i f. \text{if } f \ i \ \text{then } 1 \ \text{else } 0$ )

```

```

interpretation prob-space M
  unfolding M-def by (intro prob-space-PiM measure-pmf.prob-space-axioms)

```

```

lemma random-variable-component: random-variable (count-space UNIV) ( $\lambda f. f \ i$ )
  unfolding X-def M-def by measurable

```

```

lemma random-variable-X [measurable]: random-variable borel (X i)
  unfolding X-def M-def by measurable

```

**lemma** *distr-M-component*:  $\text{distr } M \text{ (count-space UNIV) } (\lambda f. f i) = \text{measure-pmf (bernoulli-pmf } p)$   
**proof** –  
**have**  $\text{distr } M \text{ (count-space UNIV) } (\lambda f. f i) = \text{distr } M \text{ (measure-pmf (bernoulli-pmf } p)) (\lambda f. f i)$   
**by** (rule *distr-cong*) *auto*  
**also have**  $\dots = \text{measure-pmf (bernoulli-pmf } p)$   
**unfolding** *M-def* **by** (*subst distr-PiM-component*) (*simp-all add: measure-pmf.prob-space-axioms*)  
**finally show** *?thesis* .  
**qed**

**lemma** *distr-M-X*:  
 $\text{distr } M \text{ borel } (X i) = \text{distr (measure-pmf (bernoulli-pmf } p)) \text{ borel } (\lambda b. \text{if } b \text{ then } 1 \text{ else } 0)$   
**proof** –  
**have**  $\text{distr } M \text{ borel } (X i) = \text{distr (distr } M \text{ (count-space UNIV) } (\lambda f. f i)) \text{ borel } (\lambda b. \text{if } b \text{ then } 1 \text{ else } 0 :: \text{real})$   
**by** (*subst distr-distr*) (*auto simp: M-def X-def o-def*)  
**also note** *distr-M-component[of i]*  
**finally show** *?thesis*  
**by** *simp*  
**qed**

**lemma** *X-has-expectation*:  $\text{integrable } M (X 0)$   
**proof** –  
**have**  $\text{integrable (bernoulli-pmf } p) (\lambda b. \text{if } b \text{ then } 1 \text{ else } 0 :: \text{real})$   
**by** *auto*  
**also have**  $\text{measure-pmf (bernoulli-pmf } p) = \text{distr } M \text{ (count-space UNIV) } (\lambda f. f 0)$   
**by** (*simp add: distr-M-component*)  
**also have**  $\text{integrable } \dots (\lambda b. \text{if } b \text{ then } 1 \text{ else } 0 :: \text{real}) = \text{integrable } M (X 0)$   
**unfolding** *X-def* **using** *random-variable-component* **by** (*subst integrable-distr-eq*)  
*auto*  
**finally show** *?thesis* .  
**qed**

**lemma** *indep*:  $\text{indep-vars } (\lambda-. \text{borel}) X \text{ UNIV}$   
**unfolding** *M-def X-def*  
**by** (*rule indep-vars-PiM-components'*) (*simp-all add: measure-pmf.prob-space-axioms*)

**lemma** *expectation-X*:  $\text{expectation } (X i) = p$   
**proof** –  
**have**  $\text{expectation } (X i) = \text{lebesgue-integral (distr } M \text{ (count-space UNIV) } (\lambda f. f i)) (\lambda b. \text{if } b \text{ then } 1 \text{ else } 0 :: \text{real})$   
**by** (*subst integral-distr*) (*simp-all add: random-variable-component X-def*)  
**also have**  $\text{distr } M \text{ (count-space UNIV) } (\lambda x. x i) = \text{measure-pmf (bernoulli-pmf } p)$   
**qed**

by (rule distr-M-component)  
 also have measure-pmf.expectation (bernoulli-pmf p) ( $\lambda b. \text{if } b \text{ then } 1 \text{ else } 0 :: \text{real}$ ) = p  
 using p by (subst integral-bernoulli-pmf) auto  
 finally show ?thesis .  
 qed

**theorem** AE f in M. ( $\lambda n. \text{card } \{i. i < n \wedge f i\} / n$ )  $\longrightarrow$  p

**proof** –

have AE f in M. ( $\lambda n. (\sum i < n. X i f) / \text{real } n$ )  $\longrightarrow$  expectation (X 0)

by (rule strong-law-of-large-numbers-iid)

(use indep X-has-expectation in ⟨simp-all add: distr-M-X⟩)

also have expectation (X 0) = p

by (simp add: expectation-X)

also have ( $\lambda x n. \sum i < n. X i x$ ) = ( $\lambda x n. \sum i \in \{i \in \{..<n\}. x i\}. 1$ )

by (intro ext sum.mono-neutral-cong-right) (auto simp: X-def)

also have ... = ( $\lambda x n. \text{real } (\text{card } \{i. i < n \wedge x i\})$ )

by simp

finally show ?thesis .

qed

end

end

**Acknowledgements.** I thank Sébastien Gouëzel for providing advice and context about the law of large numbers and ergodic theory. I do not actually know any ergodic theory and without him, I would probably have shied away from formalising this.

## References

- [1] Sébastien Gouëzel. Ergodic theory. *Archive of Formal Proofs*, December 2015. ISSN 2150-914x.  
[https://isa-afp.org/entries/Ergodic\\_Theory.html](https://isa-afp.org/entries/Ergodic_Theory.html), Formal proof development.
- [2] Ulrich Krengel. *Ergodic Theorems*. De Gruyter, January 1985.  
 doi:10.1515/9783110844641.
- [3] Michel Simonnet. *Measures and Probabilities*. Springer New York, New York, NY, 1996. ISBN 978-1-4612-4012-9.  
 doi:10.1007/978-1-4612-4012-9\_15.