

Kleene Algebras with Domain

Victor B. F. Gomes, Walter Guttmann, Peter Höfner,
Georg Struth and Tjark Weber

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Abstract

Kleene algebras with domain are Kleene algebras endowed with an operation that maps each element of the algebra to its domain of definition (or its complement) in abstract fashion. They form a simple algebraic basis for Hoare logics, dynamic logics or predicate transformer semantics. We formalise a modular hierarchy of algebras with domain and antidomain (domain complement) operations in Isabelle/HOL that ranges from domain and antidomain semigroups to modal Kleene algebras and divergence Kleene algebras. We link these algebras with models of binary relations and program traces. We include some examples from modal logics, termination and program analysis.

Contents

1	Introductory Remarks	2
2	Domain Semirings	3
2.1	Domain Semigroups and Domain Monoids	3
2.2	Domain Near-Semirings	5
2.3	Domain Pre-Dioids	8
2.4	Domain Semirings	9
2.5	The Algebra of Domain Elements	9
2.6	Domain Semirings with a Greatest Element	10
2.7	Forward Diamond Operators	11
2.8	Domain Kleene Algebras	12
3	Antidomain Semirings	13
3.1	Antidomain Monoids	13
3.2	Antidomain Near-Semirings	16
3.3	Antidomain Pre-Dioids	18
3.4	Antidomain Semirings	20
3.5	The Boolean Algebra of Domain Elements	21
3.6	Further Properties	22

3.7	Forward Box and Diamond Operators	24
3.8	Antidomain Kleene Algebras	26
4	Range and Antirange Semirings	27
4.1	Range Semirings	27
4.2	Antirange Semirings	28
4.3	Antirange Kleene Algebras	29
5	Modal Kleene Algebras	29
6	Models of Modal Kleene Algebras	31
7	Applications of Modal Kleene Algebras	32
7.1	A Reachability Result	32
7.2	Derivation of Segerberg’s Formula	32
7.3	Wellfoundedness and Loeb’s Formula	33
7.4	Divergence Kleene Algebras and Separation of Termination	34

1 Introductory Remarks

These theory files are intended as a reference formalisation for variants of Kleene algebras with domain. The algebraic hierarchy is developed in a modular way from domain and antidomain semigroups to modal Kleene algebras in which forward and backward box and diamond operators interact via conjugations and Galois connections. Throughout the development we have aimed at readable proofs so that these theories can be seen as a machine-checked introduction to reasoning in this setting. Apart from that, the Isabelle code is only sparsely annotated, and we refer to a series of articles for further information.

Our formalisation follows the approaches of Desharnais, Jipsen and Struth to domain semigroups [3] and Desharnais and Struth to families of domain semirings and Kleene algebras with domain [7, 6]. The link with modal Kleene algebras, Hoare logics and predicate transformers has been elaborated by Möller and Struth [13]; a notion of divergence has been added by Desharnais, Möller and Struth [5]. A previous stage of this formalisation has been documented in a companion article [11].

The target model of these axiomatisations are binary relations, where the domain operation represents the set of those elements that are related to some other element. There is a vast amount of literature on axiomatising the domain of functions, especially in semigroup theory. The deterministic nature of functions, however, leads to different axiom sets. An integration of these approaches is left for future work.

Our Isabelle/HOL formalisation itself is based on a formalisation of variants of Kleene algebras [1]. An adaptation of Kleene algebras with domain to the setting of concurrent dynamic algebra [10] can also be found in the Archive of Formal Proofs [9]. A formalisation of the original two-sorted approach to Kleene algebra with domain [4] is left for future work as well.

2 Domain Semirings

```
theory Domain-Semiring
imports Kleene-Algebra.Kleene-Algebra
```

```
begin
```

2.1 Domain Semigroups and Domain Monoids

```
class domain-op =
  fixes domain-op :: 'a  $\Rightarrow$  'a (d)
```

First we define the class of domain semigroups. Axioms are taken from [3].

```
class domain-semigroup = semigroup-mult + domain-op +
  assumes dsg1 [simp]: d x · x = x
  and dsg2 [simp]: d (x · d y) = d (x · y)
  and dsg3 [simp]: d (d x · y) = d x · d y
  and dsg4: d x · d y = d y · d x
```

```
begin
```

```
lemma domain-invol [simp]: d (d x) = d x
<proof>
```

The next lemmas show that domain elements form semilattices.

```
lemma dom-el-idem [simp]: d x · d x = d x
<proof>
```

```
lemma dom-mult-closed [simp]: d (d x · d y) = d x · d y
<proof>
```

```
lemma dom-lc3 [simp]: d x · d (x · y) = d (x · y)
<proof>
```

```
lemma d-fixpoint: ( $\exists y. x = d y$ )  $\longleftrightarrow$  x = d x
<proof>
```

```
lemma d-type:  $\forall P. (\forall x. x = d x \longrightarrow P x) \longleftrightarrow (\forall x. P (d x))$ 
<proof>
```

We define the semilattice ordering on domain semigroups and explore the semilattice of domain elements from the order point of view.

definition *ds-ord* :: 'a ⇒ 'a ⇒ bool (**infix** \sqsubseteq 50) **where**
 $x \sqsubseteq y \longleftrightarrow x = d \ x \cdot y$

lemma *ds-ord-refl*: $x \sqsubseteq x$
<proof>

lemma *ds-ord-trans*: $x \sqsubseteq y \Longrightarrow y \sqsubseteq z \Longrightarrow x \sqsubseteq z$
<proof>

lemma *ds-ord-antisym*: $x \sqsubseteq y \Longrightarrow y \sqsubseteq x \Longrightarrow x = y$
<proof>

This relation is indeed an order.

sublocale *ds*: *ordering* $\langle (\sqsubseteq) \rangle \langle \lambda x \ y. x \sqsubseteq y \wedge x \neq y \rangle$
<proof>

declare *ds.refl* [*simp*]

lemma *ds-ord-eq*: $x \sqsubseteq d \ x \longleftrightarrow x = d \ x$
<proof>

lemma $x \sqsubseteq y \Longrightarrow z \cdot x \sqsubseteq z \cdot y$

<proof>

lemma *ds-ord-iso-right*: $x \sqsubseteq y \Longrightarrow x \cdot z \sqsubseteq y \cdot z$
<proof>

The order on domain elements could as well be defined based on multiplication/meet.

lemma *ds-ord-sl-ord*: $d \ x \sqsubseteq d \ y \longleftrightarrow d \ x \cdot d \ y = d \ x$
<proof>

lemma *ds-ord-1*: $d \ (x \cdot y) \sqsubseteq d \ x$
<proof>

lemma *ds-subid-aux*: $d \ x \cdot y \sqsubseteq y$
<proof>

lemma $y \cdot d \ x \sqsubseteq y$

<proof>

lemma *ds-dom-iso*: $x \sqsubseteq y \Longrightarrow d \ x \sqsubseteq d \ y$
<proof>

lemma *ds-dom-llp*: $x \sqsubseteq d \ y \cdot x \longleftrightarrow d \ x \sqsubseteq d \ y$
<proof>

lemma *ds-dom-llp-strong*: $x = d y \cdot x \longleftrightarrow d x \sqsubseteq d y$
 ⟨*proof*⟩

definition *refines* :: 'a ⇒ 'a ⇒ bool
where *refines* $x y \equiv d y \sqsubseteq d x \wedge (d y) \cdot x \sqsubseteq y$

lemma *refines-refl*: *refines* $x x$
 ⟨*proof*⟩

lemma *refines-trans*: *refines* $x y \implies \text{refines } y z \implies \text{refines } x z$
 ⟨*proof*⟩

lemma *refines-antisym*: *refines* $x y \implies \text{refines } y x \implies x = y$
 ⟨*proof*⟩

sublocale *ref*: *ordering refines* $\lambda x y. (\text{refines } x y \wedge x \neq y)$
 ⟨*proof*⟩

end

We expand domain semigroups to domain monoids.

class *domain-monoid* = *monoid-mult* + *domain-semigroup*
begin

lemma *dom-one* [*simp*]: $d 1 = 1$
 ⟨*proof*⟩

lemma *ds-subid-eq*: $x \sqsubseteq 1 \longleftrightarrow x = d x$
 ⟨*proof*⟩

end

2.2 Domain Near-Semirings

The axioms for domain near-semirings are taken from [6].

class *domain-near-semiring* = *ab-near-semiring* + *plus-ord* + *domain-op* +
assumes *dns1* [*simp*]: $d x \cdot x = x$
and *dns2* [*simp*]: $d (x \cdot d y) = d(x \cdot y)$
and *dns3* [*simp*]: $d (x + y) = d x + d y$
and *dns4*: $d x \cdot d y = d y \cdot d x$
and *dns5* [*simp*]: $d x \cdot (d x + d y) = d x$

begin

Domain near-semirings are automatically dioids; addition is idempotent.

subclass *near-dioid*
 ⟨*proof*⟩

Next we prepare to show that domain near-semirings are domain semigroups.

lemma *dom-iso*: $x \leq y \implies d x \leq d y$
<proof>

lemma *dom-add-closed* [*simp*]: $d (d x + d y) = d x + d y$
<proof>

lemma *dom-absorp-2* [*simp*]: $d x + d x \cdot d y = d x$
<proof>

lemma *dom-1*: $d (x \cdot y) \leq d x$
<proof>

lemma *dom-subid-aux2*: $d x \cdot y \leq y$
<proof>

lemma *dom-glb*: $d x \leq d y \implies d x \leq d z \implies d x \leq d y \cdot d z$
<proof>

lemma *dom-glb-eq*: $d x \leq d y \cdot d z \iff d x \leq d y \wedge d x \leq d z$
<proof>

lemma *dom-ord*: $d x \leq d y \iff d x \cdot d y = d x$
<proof>

lemma *dom-export* [*simp*]: $d (d x \cdot y) = d x \cdot d y$
<proof>

subclass *domain-semigroup*
<proof>

We compare the domain semigroup ordering with that of the dioid.

lemma *d-two-orders*: $d x \sqsubseteq d y \iff d x \leq d y$
<proof>

lemma *two-orders*: $x \sqsubseteq y \implies x \leq y$
<proof>

lemma $x \leq y \implies x \sqsubseteq y$
<proof>

Next we prove additional properties.

lemma *dom-subdist*: $d x \leq d (x + y)$
<proof>

lemma *dom-distrib*: $d x + d y \cdot d z = (d x + d y) \cdot (d x + d z)$
<proof>

lemma *dom-llp1*: $x \leq d y \cdot x \implies d x \leq d y$

<proof>

lemma *dom-llp2*: $d x \leq d y \implies x \leq d y \cdot x$
<proof>

lemma *dom-llp*: $x \leq d y \cdot x \iff d x \leq d y$
<proof>

end

We expand domain near-semirings by an additive unit, using slightly different axioms.

class *domain-near-semiring-one* = *ab-near-semiring-one* + *plus-ord* + *domain-op*
+
assumes *dns01* [*simp*]: $x + d x \cdot x = d x \cdot x$
and *dns02* [*simp*]: $d (x \cdot d y) = d (x \cdot y)$
and *dns03* [*simp*]: $d x + 1 = 1$
and *dns04* [*simp*]: $d (x + y) = d x + d y$
and *dns05*: $d x \cdot d y = d y \cdot d x$

begin

The previous axioms are derivable.

subclass *domain-near-semiring*
<proof>

subclass *domain-monoid* *<proof>*

lemma *dom-subid*: $d x \leq 1$
<proof>

end

We add a left unit of multiplication.

class *domain-near-semiring-one-zero1* = *ab-near-semiring-one-zero1* + *domain-near-semiring-one*
+
assumes *dns06* [*simp*]: $d 0 = 0$

begin

lemma *domain-very-strict*: $d x = 0 \iff x = 0$
<proof>

lemma *dom-weakly-local*: $x \cdot y = 0 \iff x \cdot d y = 0$
<proof>

end

2.3 Domain Pre-Dioids

Pre-semirings with one and a left zero are automatically dioids. Hence there is no point defining domain pre-semirings separately from domain dioids. The axioms are once again from [6].

```

class domain-pre-dioid-one = pre-dioid-one + domain-op +
  assumes dpd1 :  $x \leq d x \cdot x$ 
  and dpd2 [simp]:  $d (x \cdot d y) = d (x \cdot y)$ 
  and dpd3 [simp]:  $d x \leq 1$ 
  and dpd4 [simp]:  $d (x + y) = d x + d y$ 

```

begin

We prepare to show that every domain pre-dioid with one is a domain near-dioid with one.

```

lemma dns1'' [simp]:  $d x \cdot x = x$ 
  <proof>

```

```

lemma d-iso:  $x \leq y \implies d x \leq d y$ 
  <proof>

```

```

lemma domain-1'':  $d (x \cdot y) \leq d x$ 
  <proof>

```

```

lemma domain-export'' [simp]:  $d (d x \cdot y) = d x \cdot d y$ 
  <proof>

```

```

lemma dom-subid-aux1'':  $d x \cdot y \leq y$ 
  <proof>

```

```

lemma dom-subid-aux2'':  $x \cdot d y \leq x$ 
  <proof>

```

```

lemma d-comm:  $d x \cdot d y = d y \cdot d x$ 
  <proof>

```

```

subclass domain-near-semiring-one
  <proof>

```

```

lemma domain-subid:  $x \leq 1 \implies x \leq d x$ 
  <proof>

```

```

lemma d-preserves-equation:  $d y \cdot x \leq x \cdot d z \iff d y \cdot x = d y \cdot x \cdot d z$ 
  <proof>

```

```

lemma d-restrict-iff:  $(x \leq y) \iff (x \leq d x \cdot y)$ 
  <proof>

```

```

lemma d-restrict-iff-1:  $(d x \cdot y \leq z) \iff (d x \cdot y \leq d x \cdot z)$ 

```


<proof>

end

We add once more a left unit of multiplication.

```
class domain-pre-doid-one-zero = domain-pre-doid-one + pre-doid-one-zero +  
  assumes dvd5 [simp]:  $d\ 0 = 0$ 
```

begin

```
subclass domain-near-semiring-one-zero  
  <proof>
```

end

2.4 Domain Semirings

We do not consider domain semirings without units separately at the moment. The axioms are taken from from [7]

```
class domain-semiringl = semiring-one-zero + plus-ord + domain-op +  
  assumes dsl1 [simp]:  $x + d\ x \cdot x = d\ x \cdot x$   
  and dsl2 [simp]:  $d\ (x \cdot d\ y) = d\ (x \cdot y)$   
  and dsl3 [simp]:  $d\ x + 1 = 1$   
  and dsl4 [simp]:  $d\ 0 = 0$   
  and dsl5 [simp]:  $d\ (x + y) = d\ x + d\ y$ 
```

begin

Every domain semiring is automatically a domain pre-doid with one and left zero.

```
subclass doid-one-zero  
  <proof>
```

```
subclass domain-pre-doid-one-zero  
  <proof>
```

end

```
class domain-semiring = domain-semiringl + semiring-one-zero
```

2.5 The Algebra of Domain Elements

We show that the domain elements of a domain semiring form a distributive lattice. Unfortunately we cannot prove this within the type class of domain semirings.

```
typedef (overloaded) 'a d-element =  $\{x :: 'a :: \text{domain-semiring}. x = d\ x\}$   
  <proof>
```

```

setup-lifting type-definition-d-element

instantiation d-element :: (domain-semiring) bounded-lattice

begin

lift-definition less-eq-d-element :: 'a d-element  $\Rightarrow$  'a d-element  $\Rightarrow$  bool is ( $\leq$ )
  <proof>

lift-definition less-d-element :: 'a d-element  $\Rightarrow$  'a d-element  $\Rightarrow$  bool is ( $<$ ) <proof>

lift-definition bot-d-element :: 'a d-element is 0
  <proof>

lift-definition top-d-element :: 'a d-element is 1
  <proof>

lift-definition inf-d-element :: 'a d-element  $\Rightarrow$  'a d-element  $\Rightarrow$  'a d-element is ( $\cdot$ )
  <proof>

lift-definition sup-d-element :: 'a d-element  $\Rightarrow$  'a d-element  $\Rightarrow$  'a d-element is
  (+)
  <proof>

instance
  <proof>

end

instance d-element :: (domain-semiring) distrib-lattice
  <proof>

```

2.6 Domain Semirings with a Greatest Element

If there is a greatest element in the semiring, then we have another equality.

```

class domain-semiring-top = domain-semiring + order-top

```

```

begin

```

```

notation top ( $\top$ )

```

```

lemma kat-equivalence-greatest:  $d\ x \leq d\ y \longleftrightarrow x \leq d\ y \cdot \top$ 
  <proof>

```

```

end

```

2.7 Forward Diamond Operators

context *domain-semiringl*

begin

We define a forward diamond operator over a domain semiring. A more modular consideration is not given at the moment.

definition $fd :: 'a \Rightarrow 'a \Rightarrow 'a ((|-) -) [61,81] 82)$ **where**
 $|x\rangle y = d (x \cdot y)$

lemma *fdia-d-simp* [*simp*]: $|x\rangle d y = |x\rangle y$
<proof>

lemma *fdia-dom* [*simp*]: $|x\rangle 1 = d x$
<proof>

lemma *fdia-add1*: $|x\rangle (y + z) = |x\rangle y + |x\rangle z$
<proof>

lemma *fdia-add2*: $|x + y\rangle z = |x\rangle z + |y\rangle z$
<proof>

lemma *fdia-mult*: $|x \cdot y\rangle z = |x\rangle |y\rangle z$
<proof>

lemma *fdia-one* [*simp*]: $|1\rangle x = d x$
<proof>

lemma *fdemodalisation1*: $d z \cdot |x\rangle y = 0 \iff d z \cdot x \cdot d y = 0$
<proof>

lemma *fdemodalisation2*: $|x\rangle y \leq d z \iff x \cdot d y \leq d z \cdot x$
<proof>

lemma *fd-iso1*: $d x \leq d y \implies |z\rangle x \leq |z\rangle y$
<proof>

lemma *fd-iso2*: $x \leq y \implies |x\rangle z \leq |y\rangle z$
<proof>

lemma *fd-zero-var* [*simp*]: $|0\rangle x = 0$
<proof>

lemma *fd-subdist-1*: $|x\rangle y \leq |x\rangle (y + z)$
<proof>

lemma *fd-subdist-2*: $|x\rangle (d y \cdot d z) \leq |x\rangle y$
<proof>

lemma *fd-subdist*: $|x\rangle (d\ y \cdot d\ z) \leq |x\rangle\ y \cdot |x\rangle\ z$
<proof>

lemma *fdia-export-1*: $d\ y \cdot |x\rangle\ z = |d\ y \cdot x\rangle\ z$
<proof>

end

context *domain-semiring*

begin

lemma *fdia-zero* [*simp*]: $|x\rangle\ 0 = 0$
<proof>

end

2.8 Domain Kleene Algebras

We add the Kleene star to our considerations. Special domain axioms are not needed.

class *domain-left-kleene-algebra* = *left-kleene-algebra-zero1* + *domain-semiring1*

begin

lemma *dom-star* [*simp*]: $d\ (x^*) = 1$
<proof>

lemma *fdia-star-unfold* [*simp*]: $|1\rangle\ y + |x\rangle\ |x^*\rangle\ y = |x^*\rangle\ y$
<proof>

lemma *fdia-star-unfoldr* [*simp*]: $|1\rangle\ y + |x^*\rangle\ |x\rangle\ y = |x^*\rangle\ y$
<proof>

lemma *fdia-star-unfold-var* [*simp*]: $d\ y + |x\rangle\ |x^*\rangle\ y = |x^*\rangle\ y$
<proof>

lemma *fdia-star-unfoldr-var* [*simp*]: $d\ y + |x^*\rangle\ |x\rangle\ y = |x^*\rangle\ y$
<proof>

lemma *fdia-star-induct-var*: $|x\rangle\ y \leq d\ y \implies |x^*\rangle\ y \leq d\ y$
<proof>

lemma *fdia-star-induct*: $d\ z + |x\rangle\ y \leq d\ y \implies |x^*\rangle\ z \leq d\ y$
<proof>

lemma *fdia-star-induct-eq*: $d\ z + |x\rangle\ y = d\ y \implies |x^*\rangle\ z \leq d\ y$
<proof>

```

end

class domain-kleene-algebra = kleene-algebra + domain-semiring

begin

subclass domain-left-kleene-algebra <proof>

end

end

```

3 Antidomain Semirings

```

theory Antidomain-Semiring
imports Domain-Semiring
begin

```

3.1 Antidomain Monoids

We axiomatise antidomain monoids, using the axioms of [3].

```

class antidomain-op =
  fixes antidomain-op :: 'a ⇒ 'a (ad)

class antidomain-left-monoid = monoid-mult + antidomain-op +
  assumes am1 [simp]: ad x · x = ad 1
  and am2: ad x · ad y = ad y · ad x
  and am3 [simp]: ad (ad x) · x = x
  and am4 [simp]: ad (x · y) · ad (x · ad y) = ad x
  and am5 [simp]: ad (x · y) · x · ad y = ad (x · y) · x

begin

```

```

no-notation domain-op (d)
no-notation zero-class.zero (0)

```

We define a zero element and operations of domain and addition.

```

definition a-zero :: 'a (0) where
  0 = ad 1

```

```

definition am-d :: 'a ⇒ 'a (d) where
  d x = ad (ad x)

```

```

definition am-add-op :: 'a ⇒ 'a ⇒ 'a (infixl ⊕ 65) where
  x ⊕ y ≡ ad (ad x · ad y)

```

```

lemma a-d-zero [simp]: ad x · d x = 0
  <proof>

```

lemma *a-d-one* [*simp*]: $d x \oplus ad x = 1$
<proof>

lemma *n-annil* [*simp*]: $0 \cdot x = 0$
<proof>

lemma *a-mult-idem* [*simp*]: $ad x \cdot ad x = ad x$
<proof>

lemma *a-add-idem* [*simp*]: $ad x \oplus ad x = ad x$
<proof>

The next three axioms suffice to show that the domain elements form a Boolean algebra.

lemma *a-add-comm*: $x \oplus y = y \oplus x$
<proof>

lemma *a-add-assoc*: $x \oplus (y \oplus z) = (x \oplus y) \oplus z$
<proof>

lemma *huntington* [*simp*]: $ad (x \oplus y) \oplus ad (x \oplus ad y) = ad x$
<proof>

lemma *a-absorb1* [*simp*]: $(ad x \oplus ad y) \cdot ad x = ad x$
<proof>

lemma *a-absorb2* [*simp*]: $ad x \oplus ad x \cdot ad y = ad x$
<proof>

The distributivity laws remain to be proved; our proofs follow those of Mad-
dux [12].

lemma *prod-split* [*simp*]: $ad x \cdot ad y \oplus ad x \cdot d y = ad x$
<proof>

lemma *sum-split* [*simp*]: $(ad x \oplus ad y) \cdot (ad x \oplus d y) = ad x$
<proof>

lemma *a-comp-simp* [*simp*]: $(ad x \oplus ad y) \cdot d x = ad y \cdot d x$
<proof>

lemma *a-distrib1*: $ad x \cdot (ad y \oplus ad z) = ad x \cdot ad y \oplus ad x \cdot ad z$
<proof>

lemma *a-distrib2*: $ad x \oplus ad y \cdot ad z = (ad x \oplus ad y) \cdot (ad x \oplus ad z)$
<proof>

lemma *aa-loc* [*simp*]: $d (x \cdot d y) = d (x \cdot y)$
<proof>

lemma *a-loc* [*simp*]: $ad (x \cdot d y) = ad (x \cdot y)$
(*proof*)

lemma *d-a-export* [*simp*]: $d (ad x \cdot y) = ad x \cdot d y$
(*proof*)

Every antidomain monoid is a domain monoid.

sublocale *dm*: *domain-monoid* *am-d* (\cdot) 1
(*proof*)

lemma *ds-ord-iso1*: $x \sqsubseteq y \implies z \cdot x \sqsubseteq z \cdot y$
(*proof*)

lemma *a-very-costrict*: $ad x = 1 \longleftrightarrow x = 0$
(*proof*)

lemma *a-weak-loc*: $x \cdot y = 0 \longleftrightarrow x \cdot d y = 0$
(*proof*)

lemma *a-closure* [*simp*]: $d (ad x) = ad x$
(*proof*)

lemma *a-d-mult-closure* [*simp*]: $d (ad x \cdot ad y) = ad x \cdot ad y$
(*proof*)

lemma *kat-3'*: $d x \cdot y \cdot ad z = 0 \implies d x \cdot y = d x \cdot y \cdot d z$
(*proof*)

lemma *s4* [*simp*]: $ad x \cdot ad (ad x \cdot y) = ad x \cdot ad y$
(*proof*)

end

class *antidomain-monoid* = *antidomain-left-monoid* +
assumes *am6* [*simp*]: $x \cdot ad 1 = ad 1$

begin

lemma *kat-3-equiv*: $d x \cdot y \cdot ad z = 0 \longleftrightarrow d x \cdot y = d x \cdot y \cdot d z$
(*proof*)

no-notation *a-zero* (0)

no-notation *am-d* (d)

end

3.2 Antidomain Near-Semirings

We define antidomain near-semirings. We do not consider units separately. The axioms are taken from [6].

notation *zero-class.zero* (0)

class *antidomain-near-semiring* = *ab-near-semiring-one-zero1* + *antidomain-op* + *plus-ord* +

assumes *ans1* [*simp*]: $ad\ x \cdot x = 0$

and *ans2* [*simp*]: $ad\ (x \cdot y) + ad\ (x \cdot ad\ (ad\ y)) = ad\ (x \cdot ad\ (ad\ y))$

and *ans3* [*simp*]: $ad\ (ad\ x) + ad\ x = 1$

and *ans4* [*simp*]: $ad\ (x + y) = ad\ x \cdot ad\ y$

begin

definition *ans-d* :: $'a \Rightarrow 'a\ (d)$ **where**
 $d\ x = ad\ (ad\ x)$

lemma *a-a-one* [*simp*]: $d\ 1 = 1$
 $\langle proof \rangle$

lemma *a-very-costrict'*: $ad\ x = 1 \iff x = 0$
 $\langle proof \rangle$

lemma *one-idem* [*simp*]: $1 + 1 = 1$
 $\langle proof \rangle$

Every antidomain near-semiring is automatically a dioid, and therefore ordered.

subclass *near-dioid-one-zero1*
 $\langle proof \rangle$

lemma *d1-a* [*simp*]: $d\ x \cdot x = x$
 $\langle proof \rangle$

lemma *a-comm*: $ad\ x \cdot ad\ y = ad\ y \cdot ad\ x$
 $\langle proof \rangle$

lemma *a-subid*: $ad\ x \leq 1$
 $\langle proof \rangle$

lemma *a-subid-aux1*: $ad\ x \cdot y \leq y$
 $\langle proof \rangle$

lemma *a-subdist*: $ad\ (x + y) \leq ad\ x$
 $\langle proof \rangle$

lemma *a-antitone*: $x \leq y \implies ad\ y \leq ad\ x$
 $\langle proof \rangle$

lemma *a-mul-d* [*simp*]: $ad\ x \cdot d\ x = 0$
 ⟨*proof*⟩

lemma *a-gla1*: $ad\ x \cdot y = 0 \implies ad\ x \leq ad\ y$
 ⟨*proof*⟩

lemma *a-gla2*: $ad\ x \leq ad\ y \implies ad\ x \cdot y = 0$
 ⟨*proof*⟩

lemma *a2-eq* [*simp*]: $ad\ (x \cdot d\ y) = ad\ (x \cdot y)$
 ⟨*proof*⟩

lemma *a-export'* [*simp*]: $ad\ (ad\ x \cdot y) = d\ x + ad\ y$
 ⟨*proof*⟩

Every antidomain near-semiring is a domain near-semiring.

sublocale *dnsz*: *domain-near-semiring-one-zero* (+) (·) 1 0 *ans-d* (≤) (<)
 ⟨*proof*⟩

lemma *a-idem* [*simp*]: $ad\ x \cdot ad\ x = ad\ x$
 ⟨*proof*⟩

lemma *a-3-var* [*simp*]: $ad\ x \cdot ad\ y \cdot (x + y) = 0$
 ⟨*proof*⟩

lemma *a-3* [*simp*]: $ad\ x \cdot ad\ y \cdot d\ (x + y) = 0$
 ⟨*proof*⟩

lemma *a-closure'* [*simp*]: $d\ (ad\ x) = ad\ x$
 ⟨*proof*⟩

The following counterexamples show that some of the antidomain monoid axioms do not need to hold.

lemma $x \cdot ad\ 1 = ad\ 1$

⟨*proof*⟩

lemma $ad\ (x \cdot y) \cdot ad\ (x \cdot ad\ y) = ad\ x$

⟨*proof*⟩

lemma $ad\ (x \cdot y) \cdot ad\ (x \cdot ad\ y) = ad\ x$

⟨*proof*⟩

lemma *phl-seq-inv*: $d\ v \cdot x \cdot y \cdot ad\ w = 0 \implies \exists z. d\ v \cdot x \cdot d\ z = 0 \wedge ad\ z \cdot y \cdot ad\ w = 0$

⟨*proof*⟩

lemma *a-fixpoint*: $ad\ x = x \implies (\forall y. y = 0)$
 ⟨*proof*⟩

no-notation *ans-d* (d)

end

3.3 Antidomain Pre-Dioids

Antidomain pre-dioids are based on a different set of axioms, which are again taken from [6].

class *antidomain-pre-dioid* = *pre-dioid-one-zero* + *antidomain-op* +
assumes *apd1* [*simp*]: $ad\ x \cdot x = 0$
and *apd2* [*simp*]: $ad\ (x \cdot y) \leq ad\ (x \cdot ad\ (ad\ y))$
and *apd3* [*simp*]: $ad\ (ad\ x) + ad\ x = 1$

begin

definition *apd-d* :: $'a \Rightarrow 'a\ (d)$ **where**
 $d\ x = ad\ (ad\ x)$

lemma *a-very-costrict'*: $ad\ x = 1 \longleftrightarrow x = 0$
 ⟨*proof*⟩

lemma *a-subid'*: $ad\ x \leq 1$
 ⟨*proof*⟩

lemma *d1-a'* [*simp*]: $d\ x \cdot x = x$
 ⟨*proof*⟩

lemma *a-subid-aux1'*: $ad\ x \cdot y \leq y$
 ⟨*proof*⟩

lemma *a-mul-d'* [*simp*]: $ad\ x \cdot d\ x = 0$
 ⟨*proof*⟩

lemma *a-d-closed* [*simp*]: $d\ (ad\ x) = ad\ x$
 ⟨*proof*⟩

lemma *meet-ord-def*: $ad\ x \leq ad\ y \longleftrightarrow ad\ x \cdot ad\ y = ad\ x$
 ⟨*proof*⟩

lemma *d-weak-loc*: $x \cdot y = 0 \longleftrightarrow x \cdot d\ y = 0$
 ⟨*proof*⟩

lemma *gla-1*: $ad\ x \cdot y = 0 \implies ad\ x \leq ad\ y$
 ⟨*proof*⟩

lemma *a2-eq'* [*simp*]: $ad (x \cdot d y) = ad (x \cdot y)$
<proof>

lemma *a-supdist-var*: $ad (x + y) \leq ad x$
<proof>

lemma *a-antitone'*: $x \leq y \implies ad y \leq ad x$
<proof>

lemma *a-comm-var*: $ad x \cdot ad y \leq ad y \cdot ad x$
<proof>

lemma *a-comm'*: $ad x \cdot ad y = ad y \cdot ad x$
<proof>

lemma *a-closed* [*simp*]: $d (ad x \cdot ad y) = ad x \cdot ad y$
<proof>

lemma *a-export''* [*simp*]: $ad (ad x \cdot y) = d x + ad y$
<proof>

lemma *d1-sum-var*: $x + y \leq (d x + d y) \cdot (x + y)$
<proof>

lemma *a4'*: $ad (x + y) = ad x \cdot ad y$
<proof>

Antidomain pre-dioids are domain pre-dioids and antidomain near-semirings,
but still not antidomain monoids.

sublocale *dpdz*: *domain-pre-diod-one-zero* (+) (·) 1 0 (≤) (<) λx. *ad* (ad x)
<proof>

subclass *antidomain-near-semiring*
<proof>

lemma *a-supdist*: $ad (x + y) \leq ad x + ad y$
<proof>

lemma *a-gla*: $ad x \cdot y = 0 \iff ad x \leq ad y$
<proof>

lemma *a-subid-aux2*: $x \cdot ad y \leq x$
<proof>

lemma *a42-var*: $d x \cdot d y \leq ad (ad x + ad y)$
<proof>

lemma *d1-weak* [*simp*]: $(d x + d y) \cdot x = x$
<proof>

lemma $x \cdot ad\ 1 = ad\ 1$

<proof>

lemma $ad\ x \cdot (y + z) = ad\ x \cdot y + ad\ x \cdot z$

<proof>

lemma $ad\ (x \cdot y) \cdot ad\ (x \cdot ad\ y) = ad\ x$

<proof>

lemma $ad\ (x \cdot y) \cdot ad\ (x \cdot ad\ y) = ad\ x$

<proof>

no-notation $apd-d\ (d)$

end

3.4 Antidomain Semirings

Antidomain semirings are direct expansions of antidomain pre-dioids, but do not require idempotency of addition. Hence we give a slightly different axiomatisation, following [7].

class *antidomain-semiring1* = *semiring-one-zero1* + *plus-ord* + *antidomain-op* +
 assumes *as1* [*simp*]: $ad\ x \cdot x = 0$
 and *as2* [*simp*]: $ad\ (x \cdot y) + ad\ (x \cdot ad\ (ad\ y)) = ad\ (x \cdot ad\ (ad\ y))$
 and *as3* [*simp*]: $ad\ (ad\ x) + ad\ x = 1$

begin

definition *ads-d* :: $'a \Rightarrow 'a\ (d)$ **where**
 $d\ x = ad\ (ad\ x)$

lemma *one-idem'*: $1 + 1 = 1$
<proof>

Every antidomain semiring is a dioid and an antidomain pre-dioid.

subclass *dioid*
<proof>

subclass *antidomain-pre-dioid*
<proof>

lemma *am5-lem* [*simp*]: $ad\ (x \cdot y) \cdot ad\ (x \cdot ad\ y) = ad\ x$
<proof>

lemma *am6-lem* [*simp*]: $ad (x \cdot y) \cdot x \cdot ad y = ad (x \cdot y) \cdot x$
 ⟨*proof*⟩

lemma *a-zero* [*simp*]: $ad 0 = 1$
 ⟨*proof*⟩

lemma *a-one* [*simp*]: $ad 1 = 0$
 ⟨*proof*⟩

subclass *antidomain-left-monoid*
 ⟨*proof*⟩

Every antidomain left semiring is a domain left semiring.

no-notation *domain-semiringl-class*.*fd* ((|-) -) [61,81] 82)

definition *fdia* :: 'a ⇒ 'a ⇒ 'a ((|-) -) [61,81] 82) **where**
 |x> y = ad (ad (x · y))

sublocale *ds*: *domain-semiringl* (+) (·) 1 0 λx. ad (ad x) (≤) (<)
rewrites *ds*.*fd* x y ≡ *fdia* x y
 ⟨*proof*⟩

lemma *fd-eq-fdia* [*simp*]: *domain-semiringl*.*fd* (·) d x y ≡ *fdia* x y
 ⟨*proof*⟩

end

class *antidomain-semiring* = *antidomain-semiringl* + *semiring-one-zero*

begin

Every antidomain semiring is an antidomain monoid.

subclass *antidomain-monoid*
 ⟨*proof*⟩

lemma *a-zero* = 0
 ⟨*proof*⟩

sublocale *ds*: *domain-semiring* (+) (·) 1 0 λx. ad (ad x) (≤) (<)
rewrites *ds*.*fd* x y ≡ *fdia* x y
 ⟨*proof*⟩

end

3.5 The Boolean Algebra of Domain Elements

typedef (**overloaded**) 'a *a2-element* = {x :: 'a :: *antidomain-semiring*. x = d x}
 ⟨*proof*⟩

```

setup-lifting type-definition-a2-element

instantiation a2-element :: (antidomain-semiring) boolean-algebra

begin

lift-definition less-eq-a2-element :: 'a a2-element  $\Rightarrow$  'a a2-element  $\Rightarrow$  bool is ( $\leq$ )
  <proof>

lift-definition less-a2-element :: 'a a2-element  $\Rightarrow$  'a a2-element  $\Rightarrow$  bool is ( $<$ )
  <proof>

lift-definition bot-a2-element :: 'a a2-element is 0
  <proof>

lift-definition top-a2-element :: 'a a2-element is 1
  <proof>

lift-definition inf-a2-element :: 'a a2-element  $\Rightarrow$  'a a2-element  $\Rightarrow$  'a a2-element
is ( $\cdot$ )
  <proof>

lift-definition sup-a2-element :: 'a a2-element  $\Rightarrow$  'a a2-element  $\Rightarrow$  'a a2-element
is ( $+$ )
  <proof>

lift-definition minus-a2-element :: 'a a2-element  $\Rightarrow$  'a a2-element  $\Rightarrow$  'a a2-element
is  $\lambda x y. x \cdot ad\ y$ 
  <proof>

lift-definition uminus-a2-element :: 'a a2-element  $\Rightarrow$  'a a2-element is antido-
main-op
  <proof>

instance
  <proof>

end

```

3.6 Further Properties

```

context antidomain-semiringl

begin

lemma a-2-var:  $ad\ x \cdot d\ y = 0 \iff ad\ x \leq ad\ y$ 
  <proof>

```

The following two lemmas give the Galois connection of Heyting algebras.

lemma *da-shunt1*: $x \leq d y + z \implies x \cdot ad y \leq z$
<proof>

lemma *da-shunt2*: $x \leq ad y + z \implies x \cdot d y \leq z$
<proof>

lemma *d-a-galois1*: $d x \cdot ad y \leq d z \iff d x \leq d z + d y$
<proof>

lemma *d-a-galois2*: $d x \cdot d y \leq d z \iff d x \leq d z + ad y$
<proof>

lemma *d-cancellation-1*: $d x \leq d y + d x \cdot ad y$
<proof>

lemma *d-cancellation-2*: $(d z + d y) \cdot ad y \leq d z$
<proof>

lemma *a-de-morgan*: $ad (ad x \cdot ad y) = d (x + y)$
<proof>

lemma *a-de-morgan-var-3*: $ad (d x + d y) = ad x \cdot ad y$
<proof>

lemma *a-de-morgan-var-4*: $ad (d x \cdot d y) = ad x + ad y$
<proof>

lemma *a-4*: $ad x \leq ad (x \cdot y)$
<proof>

lemma *a-6*: $ad (d x \cdot y) = ad x + ad y$
<proof>

lemma *a-7*: $d x \cdot ad (d y + d z) = d x \cdot ad y \cdot ad z$
<proof>

lemma *a-d-add-closure [simp]*: $d (ad x + ad y) = ad x + ad y$
<proof>

lemma *d-6 [simp]*: $d x + ad x \cdot d y = d x + d y$
<proof>

lemma *d-7 [simp]*: $ad x + d x \cdot ad y = ad x + ad y$
<proof>

lemma *a-mult-add*: $ad x \cdot (y + x) = ad x \cdot y$
<proof>

lemma *kat-2*: $y \cdot ad z \leq ad x \cdot y \implies d x \cdot y \cdot ad z = 0$

<proof>

lemma *kat-3*: $d x \cdot y \cdot ad z = 0 \implies d x \cdot y = d x \cdot y \cdot d z$
<proof>

lemma *kat-4*: $d x \cdot y = d x \cdot y \cdot d z \implies d x \cdot y \leq y \cdot d z$
<proof>

lemma *kat-2-equiv*: $y \cdot ad z \leq ad x \cdot y \iff d x \cdot y \cdot ad z = 0$
<proof>

lemma *kat-4-equiv*: $d x \cdot y = d x \cdot y \cdot d z \iff d x \cdot y \leq y \cdot d z$
<proof>

lemma *kat-3-equiv-opp*: $ad z \cdot y \cdot d x = 0 \iff y \cdot d x = d z \cdot y \cdot d x$
<proof>

lemma *kat-4-equiv-opp*: $y \cdot d x = d z \cdot y \cdot d x \iff y \cdot d x \leq d z \cdot y$
<proof>

3.7 Forward Box and Diamond Operators

lemma *fdemodalisation22*: $|x\rangle y \leq d z \iff ad z \cdot x \cdot d y = 0$
<proof>

lemma *dia-diff-var*: $|x\rangle y \leq |x\rangle (d y \cdot ad z) + |x\rangle z$
<proof>

lemma *dia-diff*: $|x\rangle y \cdot ad (|x\rangle z) \leq |x\rangle (d y \cdot ad z)$
<proof>

lemma *fdia-export-2*: $ad y \cdot |x\rangle z = |ad y \cdot x\rangle z$
<proof>

lemma *fdia-split*: $|x\rangle y = d z \cdot |x\rangle y + ad z \cdot |x\rangle y$
<proof>

definition *fbox* :: 'a \Rightarrow 'a \Rightarrow 'a ((|-|) [61,81] 82) **where**
 $|x] y = ad (x \cdot ad y)$

The next lemmas establish the De Morgan duality between boxes and diamonds.

lemma *fdia-fbox-de-morgan-2*: $ad (|x\rangle y) = |x] ad y$
<proof>

lemma *fbox-simp*: $|x] y = |x] d y$
<proof>

lemma *fbox-dom [simp]*: $|x] 0 = ad x$

<proof>

lemma *fbbox-add1*: $|x| (d y \cdot d z) = |x| y \cdot |x| z$
<proof>

lemma *fbbox-add2*: $|x + y| z = |x| z \cdot |y| z$
<proof>

lemma *fbbox-mult*: $|x \cdot y| z = |x| |y| z$
<proof>

lemma *fbbox-zero* [*simp*]: $|0| x = 1$
<proof>

lemma *fbbox-one* [*simp*]: $|1| x = d x$
<proof>

lemma *fbbox-iso*: $d x \leq d y \implies |z| x \leq |z| y$
<proof>

lemma *fbbox-antitone-var*: $x \leq y \implies |y| z \leq |x| z$
<proof>

lemma *fbbox-subdist-1*: $|x| (d y \cdot d z) \leq |x| y$
<proof>

lemma *fbbox-subdist-2*: $|x| y \leq |x| (d y + d z)$
<proof>

lemma *fbbox-subdist*: $|x| d y + |x| d z \leq |x| (d y + d z)$
<proof>

lemma *fbbox-diff-var*: $|x| (d y + ad z) \cdot |x| z \leq |x| y$
<proof>

lemma *fbbox-diff*: $|x| (d y + ad z) \leq |x| y + ad (|x| z)$
<proof>

end

context *antidomain-semiring*

begin

lemma *kat-1*: $d x \cdot y \leq y \cdot d z \implies y \cdot ad z \leq ad x \cdot y$
<proof>

lemma *kat-1-equiv*: $d x \cdot y \leq y \cdot d z \iff y \cdot ad z \leq ad x \cdot y$
<proof>

lemma *kat-3-equiv'*: $d x \cdot y \cdot ad z = 0 \iff d x \cdot y = d x \cdot y \cdot d z$
 ⟨proof⟩

lemma *kat-1-equiv-opp*: $y \cdot d x \leq d z \cdot y \iff ad z \cdot y \leq y \cdot ad x$
 ⟨proof⟩

lemma *kat-2-equiv-opp*: $ad z \cdot y \leq y \cdot ad x \iff ad z \cdot y \cdot d x = 0$
 ⟨proof⟩

lemma *fbox-one-1* [*simp*]: $|x| 1 = 1$
 ⟨proof⟩

lemma *fbox-demodalisation3*: $d y \leq |x| d z \iff d y \cdot x \leq x \cdot d z$
 ⟨proof⟩

end

3.8 Antidomain Kleene Algebras

class *antidomain-left-kleene-algebra* = *antidomain-semiringl* + *left-kleene-algebra-zero1*

begin

sublocale *dka*: *domain-left-kleene-algebra* (+) (·) 1 0 d (≤) (<) *star*
rewrites *domain-semiringl.fd* (·) $d x y \equiv |x| y$
 ⟨proof⟩

lemma *a-star* [*simp*]: $ad (x^*) = 0$
 ⟨proof⟩

lemma *fbox-star-unfold* [*simp*]: $|1| z \cdot |x| |x^*| z = |x^*| z$
 ⟨proof⟩

lemma *fbox-star-unfold-var* [*simp*]: $d z \cdot |x| |x^*| z = |x^*| z$
 ⟨proof⟩

lemma *fbox-star-unfoldr* [*simp*]: $|1| z \cdot |x^*| |x| z = |x^*| z$
 ⟨proof⟩

lemma *fbox-star-unfoldr-var* [*simp*]: $d z \cdot |x^*| |x| z = |x^*| z$
 ⟨proof⟩

lemma *fbox-star-induct-var*: $d y \leq |x| y \implies d y \leq |x^*| y$
 ⟨proof⟩

lemma *fbox-star-induct*: $d y \leq d z \cdot |x| y \implies d y \leq |x^*| z$
 ⟨proof⟩

lemma *fbox-star-induct-eq*: $d z \cdot |x] y = d y \implies d y \leq |x^*] z$
 ⟨proof⟩

lemma *fbox-export-1*: $ad y + |x] y = |d y \cdot x] y$
 ⟨proof⟩

lemma *fbox-export-2*: $d y + |x] y = |ad y \cdot x] y$
 ⟨proof⟩

end

class *antidomain-kleene-algebra* = *antidomain-semiring* + *kleene-algebra*

begin

subclass *antidomain-left-kleene-algebra* ⟨proof⟩

lemma $d p \leq |(d t \cdot x)^* \cdot ad t] (d q \cdot ad t) \implies d p \leq |d t \cdot x] d q$
 ⟨proof⟩

end

end

4 Range and Antirange Semirings

theory *Range-Semiring*
imports *Antidomain-Semiring*
begin

4.1 Range Semirings

We set up the duality between domain and antidomain semirings on the one hand and range and antirange semirings on the other hand.

class *range-op* =
fixes *range-op* :: 'a \Rightarrow 'a (r)

class *range-semiring* = *semiring-one-zero* + *plus-ord* + *range-op* +
assumes *rsr1* [*simp*]: $x + (x \cdot r x) = x \cdot r x$
and *rsr2* [*simp*]: $r (r x \cdot y) = r(x \cdot y)$
and *rsr3* [*simp*]: $r x + 1 = 1$
and *rsr4* [*simp*]: $r 0 = 0$
and *rsr5* [*simp*]: $r (x + y) = r x + r y$

begin

definition *bd* :: 'a \Rightarrow 'a \Rightarrow 'a (⟨-| - [61,81] 82) **where**

$\langle x | y = r (y \cdot x)$

no-notation *range-op* (*r*)

end

sublocale *range-semiring* \subseteq *rdual*: *domain-semiring* (+) $\lambda x y. y \cdot x$ 1 0 *range-op*
(\leq) (<)
rewrites *rdual.fd* $x y = \langle x | y$
(*proof*)

sublocale *domain-semiring* \subseteq *ddual*: *range-semiring* (+) $\lambda x y. y \cdot x$ 1 0 *domain-op*
(\leq) (<)
rewrites *ddual.bd* $x y = \text{domain-semiringl-class.fd } x y$
(*proof*)

4.2 Antirange Semirings

class *antirange-op* =
fixes *antirange-op* :: 'a \Rightarrow 'a (*ar* - [999] 1000)

class *antirange-semiring* = *semiring-one-zero* + *plus-ord* + *antirange-op* +
assumes *ars1* [*simp*]: $x \cdot \text{ar } x = 0$
and *ars2* [*simp*]: $\text{ar } (x \cdot y) + \text{ar } (\text{ar } \text{ar } x \cdot y) = \text{ar } (\text{ar } \text{ar } x \cdot y)$
and *ars3* [*simp*]: $\text{ar } \text{ar } x + \text{ar } x = 1$

begin

no-notation *bd* ($\langle - | -$ - [61,81] 82)

definition *ars-r* :: 'a \Rightarrow 'a (*r*) **where**
 $r x = \text{ar } (\text{ar } x)$

definition *bdia* :: 'a \Rightarrow 'a \Rightarrow 'a ($\langle - | -$ - [61,81] 82) **where**
 $\langle x | y = \text{ar } \text{ar } (y \cdot x)$

definition *bbox* :: 'a \Rightarrow 'a \Rightarrow 'a ($[- | -$ - [61,81] 82) **where**
 $[x | y = \text{ar } (\text{ar } y \cdot x)$

end

sublocale *antirange-semiring* \subseteq *ardual*: *antidomain-semiring* *antirange-op* (+) λx
 $y. y \cdot x$ 1 0 (\leq) (<)
rewrites *ardual.ads-d* $x = r x$
and *ardual.fdia* $x y = \langle x | y$
and *ardual.fbox* $x y = [x | y$
(*proof*)

context *antirange-semiring*

begin

sublocale *rs*: *range-semiring* (+) (·) 1 0 $\lambda x. ar\ ar\ x$ (\leq) ($<$)
⟨*proof*⟩

end

sublocale *antidomain-semiring* \subseteq *addual*: *antirange-semiring* (+) $\lambda x\ y. y \cdot x$ 1 0
antidomain-op (\leq) ($<$)
rewrites *addual.ars-r* $x = d\ x$
and *addual.bdia* $x\ y = |x\rangle\ y$
and *addual.bbox* $x\ y = |x]\ y$
⟨*proof*⟩

4.3 Antirange Kleene Algebras

class *antirange-kleene-algebra* = *antirange-semiring* + *kleene-algebra*

sublocale *antirange-kleene-algebra* \subseteq *dual*: *antidomain-kleene-algebra* *antirange-op*
(+) $\lambda x\ y. y \cdot x$ 1 0 (\leq) ($<$) *star*
⟨*proof*⟩

sublocale *antidomain-kleene-algebra* \subseteq *dual*: *antirange-kleene-algebra* (+) $\lambda x\ y. y$
 $\cdot\ x$ 1 0 (\leq) ($<$) *star* *antidomain-op*
⟨*proof*⟩

Hence all range theorems have been derived by duality in a generic way.

end

5 Modal Kleene Algebras

This section studies laws that relate antidomain and antirange semirings and Kleene algebra, notably Galois connections and conjugations as those studied in [13, 7].

theory *Modal-Kleene-Algebra*
imports *Range-Semiring*
begin

class *modal-semiring* = *antidomain-semiring* + *antirange-semiring* +
assumes *domrange* [*simp*]: $d\ (r\ x) = r\ x$
and *rangedom* [*simp*]: $r\ (d\ x) = d\ x$

begin

These axioms force that the domain algebra and the range algebra coincide.

lemma *domrangefix*: $d\ x = x \longleftrightarrow r\ x = x$

<proof>

lemma *box-diamond-galois-1*:
assumes $d\ p = p$ **and** $d\ q = q$
shows $\langle x \mid p \leq q \longleftrightarrow p \leq [x]\ q$
<proof>

lemma *box-diamond-galois-2*:
assumes $d\ p = p$ **and** $d\ q = q$
shows $[x]\ p \leq q \longleftrightarrow p \leq [x]\ q$
<proof>

lemma *diamond-conjugation-var-1*:
assumes $d\ p = p$ **and** $d\ q = q$
shows $[x]\ p \leq ad\ q \longleftrightarrow \langle x \mid q \leq ad\ p$
<proof>

lemma *diamond-conjugation-aux*:
assumes $d\ p = p$ **and** $d\ q = q$
shows $\langle x \mid p \leq ad\ q \longleftrightarrow q \cdot \langle x \mid p = 0$
<proof>

lemma *diamond-conjugation*:
assumes $d\ p = p$ **and** $d\ q = q$
shows $p \cdot [x]\ q = 0 \longleftrightarrow q \cdot \langle x \mid p = 0$
<proof>

lemma *box-conjugation-var-1*:
assumes $d\ p = p$ **and** $d\ q = q$
shows $ad\ p \leq [x]\ q \longleftrightarrow ad\ q \leq [x]\ p$
<proof>

lemma *box-diamond-cancellation-1*: $d\ p = p \implies p \leq [x]\ \langle x \mid p$
<proof>

lemma *box-diamond-cancellation-2*: $d\ p = p \implies p \leq [x]\ [x]\ p$
<proof>

lemma *box-diamond-cancellation-3*: $d\ p = p \implies [x]\ [x]\ p \leq p$
<proof>

lemma *box-diamond-cancellation-4*: $d\ p = p \implies \langle x \mid [x]\ p \leq p$
<proof>

end

class *modal-kleene-algebra* = *modal-semiring* + *kleene-algebra*
begin

```

subclass antidomain-kleene-algebra <proof>

subclass antirange-kleene-algebra <proof>

end

end

```

6 Models of Modal Kleene Algebras

```

theory Modal-Kleene-Algebra-Models
imports Kleene-Algebra.Kleene-Algebra-Models
         Modal-Kleene-Algebra

```

```

begin

```

This section develops the relation model. We also briefly develop the trace model for antidomain Kleene algebras, but not for antirange or full modal Kleene algebras. The reason is that traces are implemented as lists; we therefore expect tedious inductive proofs in the presence of range. The language model is not particularly interesting.

```

definition rel-ad :: 'a rel  $\Rightarrow$  'a rel where
  rel-ad R = {(x,x) | x.  $\neg$  ( $\exists$  y. (x,y)  $\in$  R)}

```

```

interpretation rel-antidomain-kleene-algebra: antidomain-kleene-algebra rel-ad ( $\cup$ )
(O) Id {} ( $\subseteq$ ) ( $\subset$ ) rtrancl
<proof>

```

```

definition trace-a :: ('p, 'a) trace set  $\Rightarrow$  ('p, 'a) trace set where
  trace-a X = {(p,[]) | p.  $\neg$  ( $\exists$  x. x  $\in$  X  $\wedge$  p = first x)}

```

```

interpretation trace-antidomain-kleene-algebra: antidomain-kleene-algebra trace-a
( $\cup$ ) t-prod t-one t-zero ( $\subseteq$ ) ( $\subset$ ) t-star
<proof>

```

The trace model should be extended to cover modal Kleene algebras in the future.

```

definition rel-ar :: 'a rel  $\Rightarrow$  'a rel where
  rel-ar R = {(y,y) | y.  $\neg$  ( $\exists$  x. (x,y)  $\in$  R)}

```

```

interpretation rel-antirange-kleene-algebra: antirange-semiring ( $\cup$ ) (O) Id {} rel-ar
( $\subseteq$ ) ( $\subset$ )
<proof>

```

```

interpretation rel-modal-kleene-algebra: modal-kleene-algebra ( $\cup$ ) (O) Id {} ( $\subseteq$ )
( $\subset$ ) rtrancl rel-ad rel-ar
<proof>

```

end

7 Applications of Modal Kleene Algebras

theory *Modal-Kleene-Algebra-Applications*
imports *Antidomain-Semiring*
begin

This file collects some applications of the theories developed so far. These are described in [11].

context *antidomain-kleene-algebra*
begin

7.1 A Reachability Result

This example is taken from [4].

lemma *opti-iterate-var* [simp]: $|(ad\ y \cdot x)^*| y = |x^*| y$
<proof>

lemma *opti-iterate* [simp]: $d\ y + |(x \cdot ad\ y)^*| |x| y = |x^*| y$
<proof>

lemma *opti-iterate-var-2* [simp]: $d\ y + |ad\ y \cdot x| |x^*| y = |x^*| y$
<proof>

7.2 Derivation of Segerberg's Formula

This example is taken from [5].

definition *Alpha* :: $'a \Rightarrow 'a \Rightarrow 'a$ (*A*)
where $A\ x\ y = d\ (x \cdot y) \cdot ad\ y$

lemma *A-dom* [simp]: $d\ (A\ x\ y) = A\ x\ y$
<proof>

lemma *A-fdia*: $A\ x\ y = |x| y \cdot ad\ y$
<proof>

lemma *A-fdia-var*: $A\ x\ y = |x| d\ y \cdot ad\ y$
<proof>

lemma *a-A*: $ad\ (A\ x\ (ad\ y)) = |x| y + ad\ y$
<proof>

lemma *fsegerberg* [simp]: $d\ y + |x^*| A\ x\ y = |x^*| y$
<proof>

lemma *fbox-segerberg* [simp]: $d\ y \cdot |x^*| (|x| y + ad\ y) = |x^*| y$
<proof>

7.3 Wellfoundedness and Loeb's Formula

This example is taken from [7].

definition *Omega* :: 'a \Rightarrow 'a \Rightarrow 'a (Ω)
where $\Omega\ x\ y = d\ y \cdot ad\ (x \cdot y)$

If y is a set, then $\Omega(x, y)$ describes those elements in y from which no further x transitions are possible.

lemma *omega-fdia*: $\Omega\ x\ y = d\ y \cdot ad\ (|x\rangle\ y)$
 $\langle proof \rangle$

lemma *omega-fbox*: $\Omega\ x\ y = d\ y \cdot |x| (ad\ y)$
 $\langle proof \rangle$

lemma *omega-absorb1* [*simp*]: $\Omega\ x\ y \cdot ad\ (|x\rangle\ y) = \Omega\ x\ y$
 $\langle proof \rangle$

lemma *omega-absorb2* [*simp*]: $\Omega\ x\ y \cdot ad\ (x \cdot y) = \Omega\ x\ y$
 $\langle proof \rangle$

lemma *omega-le-1*: $\Omega\ x\ y \leq d\ y$
 $\langle proof \rangle$

lemma *omega-subid*: $\Omega\ x\ (d\ y) \leq d\ y$
 $\langle proof \rangle$

lemma *omega-le-2*: $\Omega\ x\ y \leq ad\ (|x\rangle\ y)$
 $\langle proof \rangle$

lemma *omega-dom* [*simp*]: $d\ (\Omega\ x\ y) = \Omega\ x\ y$
 $\langle proof \rangle$

lemma *a-omega*: $ad\ (\Omega\ x\ y) = ad\ y + |x\rangle\ y$
 $\langle proof \rangle$

lemma *omega-fdia-3* [*simp*]: $d\ y \cdot ad\ (\Omega\ x\ y) = d\ y \cdot |x\rangle\ y$
 $\langle proof \rangle$

lemma *omega-zero-equiv-1*: $\Omega\ x\ y = 0 \iff d\ y \leq |x\rangle\ y$
 $\langle proof \rangle$

definition *Loebian* :: 'a \Rightarrow bool
where *Loebian* $x = (\forall y. |x\rangle\ y \leq |x\rangle\ \Omega\ x\ y)$

definition *PreLoebian* :: 'a \Rightarrow bool
where *PreLoebian* $x = (\forall y. d\ y \leq |x^*\rangle\ \Omega\ x\ y)$

definition *Noetherian* :: 'a \Rightarrow bool
where *Noetherian* $x = (\forall y. \Omega\ x\ y = 0 \implies d\ y = 0)$

lemma *noetherian-alt*: $\text{Noetherian } x \longleftrightarrow (\forall y. d \ y \leq |x\rangle y \longrightarrow d \ y = 0)$
 ⟨proof⟩

lemma *Noetherian-iff-PreLoebian*: $\text{Noetherian } x \longleftrightarrow \text{PreLoebian } x$
 ⟨proof⟩

lemma *Loebian-imp-Noetherian*: $\text{Loebian } x \implies \text{Noetherian } x$
 ⟨proof⟩

lemma *d-transitive*: $(\forall y. |x\rangle |x\rangle y \leq |x\rangle y) \implies (\forall y. |x\rangle y = |x^*\rangle |x\rangle y)$
 ⟨proof⟩

lemma *d-transitive-var*: $(\forall y. |x\rangle |x\rangle y \leq |x\rangle y) \implies (\forall y. |x\rangle y = |x\rangle |x^*\rangle y)$
 ⟨proof⟩

lemma *d-transitive-PreLoebian-imp-Loebian*: $(\forall y. |x\rangle |x\rangle y \leq |x\rangle y) \implies \text{PreLoebian } x \implies \text{Loebian } x$
 ⟨proof⟩

lemma *d-transitive-Noetherian-iff-Loebian*: $\forall y. |x\rangle |x\rangle y \leq |x\rangle y \implies \text{Noetherian } x \longleftrightarrow \text{Loebian } x$
 ⟨proof⟩

lemma *Loeb-iff-box-Loeb*: $\text{Loebian } x \longleftrightarrow (\forall y. |x\rangle (ad \ (|x\rangle y) + d \ y) \leq |x\rangle y)$
 ⟨proof⟩

end

7.4 Divergence Kleene Algebras and Separation of Termination

The notion of divergence has been added to modal Kleene algebras in [5]. More facts about divergence could be added in the future. Some could be adapted from omega algebras.

```

class nabla-op =
  fixes nabla :: 'a  $\Rightarrow$  'a ( $\nabla$ - [999] 1000)

class fdivergence-kleene-algebra = antidomain-kleene-algebra + nabla-op +
  assumes nabla-closure [simp]:  $d \ \nabla \ x = \nabla \ x$ 
  and nabla-unfold:  $\nabla \ x \leq |x\rangle \nabla \ x$ 
  and nabla-coinduction:  $d \ y \leq |x\rangle y + d \ z \implies d \ y \leq \nabla \ x + |x^*\rangle z$ 

begin

lemma nabla-coinduction-var:  $d \ y \leq |x\rangle y \implies d \ y \leq \nabla \ x$ 
  ⟨proof⟩

```

lemma *nabla-unfold-eq* [simp]: $|x\rangle \nabla x = \nabla x$
<proof>

lemma *nabla-subdist*: $\nabla x \leq \nabla (x + y)$
<proof>

lemma *nabla-iso*: $x \leq y \implies \nabla x \leq \nabla y$
<proof>

lemma *nabla-omega*: $\Omega x (d y) = 0 \implies d y \leq \nabla x$
<proof>

lemma *nabla-noether*: $\nabla x = 0 \implies \text{Noetherian } x$
<proof>

lemma *nabla-preloeb*: $\nabla x = 0 \implies \text{PreLoebian } x$
<proof>

lemma *star-nabla-1* [simp]: $|x^*\rangle \nabla x = \nabla x$
<proof>

lemma *nabla-sum-expand* [simp]: $|x\rangle \nabla (x + y) + |y\rangle \nabla (x + y) = \nabla (x + y)$
<proof>

lemma *wagner-3*:
 assumes $d z + |x\rangle \nabla (x + y) = \nabla (x + y)$
 shows $\nabla (x + y) = \nabla x + |x^*\rangle z$
<proof>

lemma *nabla-sum-unfold* [simp]: $\nabla x + |x^*\rangle |y\rangle \nabla (x + y) = \nabla (x + y)$
<proof>

lemma *nabla-separation*: $y \cdot x \leq x \cdot (x + y)^* \implies (\nabla (x + y) = \nabla x + |x^*\rangle \nabla y)$
<proof>

The next lemma is a separation of termination theorem by Bachmair and Dershowitz [2].

lemma *bachmair-dershowitz*: $y \cdot x \leq x \cdot (x + y)^* \implies \nabla x + \nabla y = 0 \iff \nabla (x + y) = 0$
<proof>

The next lemma is a more complex separation of termination theorem by Doornbos, Backhouse and van der Woude [8].

lemma *separation-of-termination*:
 assumes $y \cdot x \leq x \cdot (x + y)^* + y$
 shows $\nabla x + \nabla y = 0 \iff \nabla (x + y) = 0$
<proof>

The final examples can be found in [11].

definition $T :: 'a \Rightarrow 'a \Rightarrow 'a \Rightarrow 'a \ (- \rightsquigarrow - \rightsquigarrow - [61,61,61] 60)$
where $p \rightsquigarrow x \rightsquigarrow q \equiv ad\ p + |x| \ d\ q$

lemma $T-d$ [simp]: $d\ (p \rightsquigarrow x \rightsquigarrow q) = p \rightsquigarrow x \rightsquigarrow q$
 ⟨proof⟩

lemma $T-p$: $d\ p \cdot (p \rightsquigarrow x \rightsquigarrow q) = d\ p \cdot |x| \ d\ q$
 ⟨proof⟩

lemma $T-a$ [simp]: $ad\ p \cdot (p \rightsquigarrow x \rightsquigarrow q) = ad\ p$
 ⟨proof⟩

lemma $T-seq$: $(p \rightsquigarrow x \rightsquigarrow q) \cdot |x|(q \rightsquigarrow y \rightsquigarrow s) \leq p \rightsquigarrow x \cdot y \rightsquigarrow s$
 ⟨proof⟩

lemma $T-square$: $(p \rightsquigarrow x \rightsquigarrow q) \cdot |x|(q \rightsquigarrow x \rightsquigarrow p) \leq p \rightsquigarrow x \cdot x \rightsquigarrow p$
 ⟨proof⟩

lemma $T-segerberg$ [simp]: $d\ p \cdot |x^*|(p \rightsquigarrow x \rightsquigarrow p) = |x^*| \ d\ p$
 ⟨proof⟩

lemma $lookahead$ [simp]: $|x^*|(d\ p \cdot |x| \ d\ p) = |x^*| \ d\ p$
 ⟨proof⟩

lemma $alternation$: $d\ p \cdot |x^*|((p \rightsquigarrow x \rightsquigarrow q) \cdot (q \rightsquigarrow x \rightsquigarrow p)) = |(x \cdot x)^*|(d\ p \cdot (q \rightsquigarrow x \rightsquigarrow p))$
 $\cdot |x \cdot (x \cdot x)^*|(d\ q \cdot (p \rightsquigarrow x \rightsquigarrow q))$
 ⟨proof⟩

lemma $|(x \cdot x)^*| \ d\ p \cdot |x \cdot (x \cdot x)^*| \ ad\ p = d\ p \cdot |x^*|((p \rightsquigarrow x \rightsquigarrow ad\ p) \cdot (ad\ p \rightsquigarrow x \rightsquigarrow p))$
 ⟨proof⟩

lemma $|x^*| \ d\ p = d\ p \cdot |x^*|(p \rightsquigarrow x \rightsquigarrow p)$
 ⟨proof⟩

end

end

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