

# Irrational numbers from THE BOOK

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## Abstract

An elementary proof is formalised: that  $\exp r$  is irrational for every nonzero rational number  $r$ . The mathematical development comes from the well-known volume *Proofs from THE BOOK* [1, pp. 51–2], by Aigner and Ziegler, who credit the idea to Hermite. The development illustrates a number of basic Isabelle techniques: the manipulation of summations, the calculation of quite complicated derivatives and the estimation of integrals. We also see how to import another AFP entry (Stirling’s formula) [2].

As for the theorem itself, note that a much stronger and more general result (the Hermite–Lindemann–Weierstraß transcendence theorem) is already available in the AFP [3].

## Contents

<b>1</b>	<b>Some irrational numbers</b>	<b>3</b>
1.1	Library additions . . . . .	3
1.2	Basic definitions and their consequences . . . . .	3
1.3	Towards the main result . . . . .	6

# 1 Some irrational numbers

From Aigner and Ziegler, *Proofs from THE BOOK* (Springer, 2018), Chapter 8, pp. 50–51.

**theory** *Irrationals-From-THEBOOK* **imports** *Stirling-Formula.Stirling-Formula*

**begin**

## 1.1 Library additions

**context** *comm-monoid-set*

**begin**

**lemma** *atLeast-atMost-pred-shift*:

$F (g \circ (\lambda n. n - \text{Suc } 0)) \{ \text{Suc } m.. \text{Suc } n \} = F g \{ m..n \}$

**unfolding** *atLeast-Suc-atMost-Suc-shift* **by** *simp*

**end**

**lemma** *field-differentiable-diff-const* [*simp, derivative-intros*]:

$(-)_c$  *field-differentiable*  $F$

**unfolding** *field-differentiable-def*

**by** (*rule derivative-eq-intros exI | force*)**+**

## 1.2 Basic definitions and their consequences

**definition** *hf* **where**  $hf \equiv \lambda x. \lambda n. \lambda x::\text{real}. (x^n * (1-x)^n) / \text{fact } n$

**definition** *cf* **where**  $cf \equiv \lambda n i. \text{if } i < n \text{ then } 0 \text{ else } (n \text{ choose } (i-n)) * (-1)^{i-n}$

Mere knowledge that the coefficients are integers is not enough later on.

**lemma** *hf-int-poly*:

**fixes**  $x::\text{real}$

**shows**  $hf\ n = (\lambda x. (1 / \text{fact } n) * (\sum_{i=0..2*n}. \text{real-of-int } (cf\ n\ i) * x^i))$

**proof**

**fix**  $x$

**have** *inj*: *inj-on*  $((+)_n) \{..n\}$

**by** (*auto simp: inj-on-def*)

**have** [*simp*]:  $((+)_n) ' \{..n\} = \{n..2*n\}$

**using** *nat-le-iff-add* **by** *fastforce*

**have**  $(x^n * (-x + 1)^n) = x^n * (\sum_{k \leq n}. \text{real } (n \text{ choose } k) * (-x)^k)$

**unfolding** *binomial-ring* **by** *simp*

**also have**  $\dots = x^n * (\sum_{k \leq n}. \text{real-of-int } ((n \text{ choose } k) * (-1)^k) * x^{n+k})$

**by** (*simp add: mult.assoc flip: power-minus*)

**also have**  $\dots = (\sum_{k \leq n}. \text{real-of-int } ((n \text{ choose } k) * (-1)^k) * x^{n+k})$

**by** (*simp add: sum-distrib-left mult-ac power-add*)

**also have**  $\dots = (\sum_{i=n..2*n}. \text{real-of-int } (cf\ n\ i) * x^i)$

**by** (*simp add: sum.reindex [OF inj, simplified] cf-def*)

**finally have**  $hf\ n\ x = (1 / \text{fact } n) * (\sum_{i = n..2 * n}. \text{real-of-int } (cf\ n\ i) * x^i)$

by (simp add: hf-def)  
 moreover have  $(\sum i = 0..<n. \text{real-of-int } (cf\ n\ i) * x^{\widehat{i}}) = 0$   
 by (simp add: cf-def)  
 ultimately show  $hf\ n\ x = (1 / \text{fact } n) * (\sum i = 0..2 * n. \text{real-of-int } (cf\ n\ i) * x^{\widehat{i}})$   
 using sum.union-disjoint [of  $\{0..<n\}$   $\{n..2*n\}$   $\lambda i. \text{real-of-int } (cf\ n\ i) * x^{\widehat{i}}$ ]  
 by (simp add: ivl-disj-int-two(7) ivl-disj-un-two(7) mult-2)  
 qed

Lemma (ii) in the text has strict inequalities, but it takes more work and is less useful.

**lemma**

assumes  $0 \leq x$   
 shows *hf-nonneg*:  $0 \leq hf\ n\ x$  and *hf-le-inverse-fact*:  $hf\ n\ x \leq 1 / \text{fact } n$   
 using *assms* by (auto simp: hf-def divide-simps mult-le-one power-le-one)

**lemma** *hf-differt* [iff]: *hf n differentiable at x*

unfolding *hf-int-poly differentiable-def*  
 by (intro derivative-eq-intros exI | simp)+

**lemma** *deriv-sum-int*:

$deriv\ (\lambda x. \sum i=0..n. \text{real-of-int } (c\ i) * x^{\widehat{i}})\ x$   
 $= (\text{if } n=0 \text{ then } 0 \text{ else } (\sum i=0..n - \text{Suc } 0. \text{real-of-int } ((\text{int } i + 1) * c\ (\text{Suc } i)) * x^{\widehat{i}}))$   
 (is *deriv ?f x = (if n=0 then 0 else ?g)*)

**proof** –

have (*?f has-real-derivative ?g*) (at  $x$ ) if  $n > 0$

**proof** –

have  $(\sum i = 0..n. i * x^{\widehat{(i - \text{Suc } 0)}} * (c\ i))$   
 $= (\sum i = \text{Suc } 0..n. (\text{real } (i - \text{Suc } 0) + 1) * \text{real-of-int } (c\ i) * x^{\widehat{(i - \text{Suc } 0)}})$

using *that* by (auto simp add: sum.atLeast-Suc-atMost intro!: sum.cong)

also have  $\dots = \text{sum } ((\lambda i. (\text{real } i + 1) * \text{real-of-int } (c\ (\text{Suc } i)) * x^{\widehat{i}}) \circ (\lambda n. n - \text{Suc } 0)) \{ \text{Suc } 0.. \text{Suc } (n - \text{Suc } 0) \}$

using *that* by *simp*

also have  $\dots = ?g$

by (*simp flip: sum.atLeast-atMost-pred-shift* [where  $m=0$ ])

finally have  $\S: (\sum a = 0..n. a * x^{\widehat{(a - \text{Suc } 0)}} * (c\ a)) = ?g$ .

show *?thesis*

by (*rule derivative-eq-intros*  $\S$  | *simp*)+

**qed**

then show *?thesis*

by (*force intro: DERIV-imp-deriv*)

**qed**

We calculate the coefficients of the  $k$ th derivative precisely.

**lemma** *hf-deriv-int-poly*:

$(deriv\ \widetilde{k})\ (hf\ n) = (\lambda x. (1 / \text{fact } n) * (\sum i=0..2*n-k. \text{real-of-int } (\text{int}(\prod \{i <..i+k\}) * cf\ n\ (i+k)) * x^{\widehat{i}})$

```

proof (induction k)
  case 0
  show ?case
    by (simp add: hf-int-poly)
next
  case (Suc k)
  define F where F  $\equiv \lambda x. (\sum i = 0..2*n - k. \text{real-of-int } (\text{int}(\prod \{i<..i+k\}) * \text{cf } n (i+k)) * x^i)$ 
  have Fd: F field-differentiable at x for x
    unfolding field-differentiable-def F-def
    by (rule derivative-eq-intros exI | force)+
  have [simp]: prod int {i<..Suc (i + k)} = (1 + int i) * prod int {Suc i<..Suc (i + k)} for i
    by (metis Suc-le-mono atLeastSucAtMost-greaterThanAtMost le-add1 of-nat-Suc prod.head)
  have deriv ( $\lambda x. F x / \text{fact } n$ ) x
    = ( $\sum i = 0..2 * n - \text{Suc } k. \text{real-of-int } (\text{int}(\prod \{i<..i+ \text{Suc } k\}) * \text{cf } n (\text{Suc } (i+k))) * x^i$ ) / fact n for x
    unfolding deriv-cdivide-right [OF Fd]
    by (fastforce simp add: F-def deriv-sum-int cf-def simp flip: of-int-mult intro: sum.cong)
  then show ?case
    by (simp add: Suc F-def)
qed

```

```

lemma hf-deriv-0: (deriv  $\sim^k$ ) (hf n) 0  $\in \mathbb{Z}$ 
proof (cases n  $\leq$  k)
  case True
  then obtain j where (fact k::real) = real-of-int j * fact n
    using fact-dvd
  by (metis dvd-def fact-nonzero mult.commute nonzero-mult-div-cancel-left of-int-fact real-of-int-div)
  moreover have prod real {0<..k} = fact k
    by (simp add: fact-prod atLeastSucAtMost-greaterThanAtMost)
  ultimately show ?thesis
    by (simp add: hf-deriv-int-poly dvd-def)
next
  case False
  then show ?thesis
    by (simp add: hf-deriv-int-poly cf-def)
qed

```

```

lemma deriv-hf-minus: deriv (hf n) = ( $\lambda x. - \text{deriv } (hf n) (1-x)$ )
proof
  fix x
  have hf n = hf n  $\circ (\lambda x. (1-x))$ 
    by (simp add: fun-eq-iff hf-def mult.commute)
  then have deriv (hf n) x = deriv (hf n  $\circ (\lambda x. (1-x))$ ) x
    by fastforce

```

**also have** ... =  $\text{deriv } (hf \ n) \ (1-x) * \text{deriv } ((-)\ 1) \ x$   
**by** (*intro real-derivative-chain*) *auto*  
**finally show**  $\text{deriv } (hf \ n) \ x = - \text{deriv } (hf \ n) \ (1-x)$   
**by** *simp*  
**qed**

**lemma** *deriv-n-hf-diff* [*iff*]:  $(\text{deriv} \ \hat{\sim} \ k) \ (hf \ n)$  *field-differentiable at x*  
**unfolding** *field-differentiable-def hf-deriv-int-poly*  
**by** (*rule derivative-eq-intros exI | force*)+

**lemma** *deriv-n-hf-minus*:  $(\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) = (\lambda x. \ (-1) \ \hat{\sim} \ k * (\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \ (1-x))$

**proof** (*induction k*)

**case** *0*

**then show** *?case*

**by** (*simp add: fun-eq-iff hf-def*)

**next**

**case** (*Suc k*)

**have** *o*:  $(\lambda x. \ (\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \ (1-x)) = (\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \circ \ (-)\ 1$

**by** *auto*

**show** *?case*

**proof**

**fix** *x*

**have** [*simp*]:  $((\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \circ \ (-)\ 1)$  *field-differentiable at x*

**by** (*force intro: field-differentiable-compose*)

**have**  $(\text{deriv} \ \hat{\sim} \ \text{Suc } k) \ (hf \ n) \ x = \text{deriv } (\lambda x. \ (-1) \ \hat{\sim} \ k * (\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \ (1-x)) \ x$

**by** *simp (metis Suc)*

**also have** ... =  $(-1) \ \hat{\sim} \ k * \text{deriv } (\lambda x. \ (\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \ (1-x)) \ x$

**using** *o* **by** *fastforce*

**also have** ... =  $(-1) \ \hat{\sim} \ \text{Suc } k * (\text{deriv} \ \hat{\sim} \ \text{Suc } k) \ (hf \ n) \ (1-x)$

**by** (*subst o, subst deriv-chain, auto*)

**finally show**  $(\text{deriv} \ \hat{\sim} \ \text{Suc } k) \ (hf \ n) \ x = (-1) \ \hat{\sim} \ \text{Suc } k * (\text{deriv} \ \hat{\sim} \ \text{Suc } k) \ (hf \ n) \ (1-x)$ .

**qed**

**qed**

### 1.3 Towards the main result

**lemma** *hf-deriv-1*:  $(\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) \ 1 \in \mathbb{Z}$

**by** (*smt (verit, best) Ints-1 Ints-minus Ints-mult Ints-power deriv-n-hf-minus hf-deriv-0*)

**lemma** *hf-deriv-eq-0*:  $k > 2*n \implies (\text{deriv} \ \hat{\sim} \ k) \ (hf \ n) = (\lambda x. \ 0)$

**by** (*force simp add: cf-def hf-deriv-int-poly*)

The case for positive integers

**lemma** *exp-nat-irrational*:

**assumes**  $s > 0$  **shows**  $\text{exp } (\text{real-of-int } s) \notin \mathbb{Q}$

**proof**

```

assume exp (real-of-int s) ∈ Q
then obtain a b where ab:  $a > 0$   $b > 0$  coprime a b and exp-s: exp s = of-int
a / of-int b
  using Rats-cases' div-0 exp-not-eq-zero of-int-0
  by (smt (verit, best) exp-gt-zero of-int-0-less-iff zero-less-divide-iff)
define n where  $n \equiv \text{nat } (\max (a^2) (3 * s^3))$ 
then have ns3:  $s^3 \leq \text{real } n / 3$ 
  by linarith
have  $n > 0$ 
  using  $\langle a > 0 \rangle$  n-def by (smt (verit, best) zero-less-nat-eq zero-less-power)
then have  $s^{(2*n+1)} \leq s^{(3*n)}$ 
  using  $\langle a > 0 \rangle$  assms by (intro power-increasing) auto
also have ... = real-of-int( $s^3$ )n
  by (simp add: power-mult)
also have ... ≤  $(n / 3)^n$ 
  using assms ns3 by (simp add: power-mono)
also have ... ≤  $(n / \text{exp } 1)^n$ 
  using exp-le  $\langle n > 0 \rangle$ 
  by (auto simp add: divide-simps)
finally have s-le:  $s^{(2*n+1)} \leq (n / \text{exp } 1)^n$ 
  by presburger
have a-less:  $a < \text{sqrt } (2*\text{pi}*n)$ 
proof –
  have  $2*\text{pi} > 1$ 
  by (smt (z3) pi-gt-zero sin-gt-zero-02 sin-le-zero)
have  $a = \text{sqrt } (a^2)$ 
  by (simp add: ab(1) order-less-imp-le)
also have ... ≤  $\text{sqrt } n$ 
  unfolding n-def
  by (smt (verit, ccfv-SIG) int-nat-eq of-nat-less-of-int-iff real-sqrt-le-mono)
also have ... <  $\text{sqrt } (2*\text{pi}*n)$ 
  by (simp add:  $\langle 0 < n \rangle$   $\langle 1 < 2 * \text{pi} \rangle$ )
finally show ?thesis .
qed
have  $\text{sqrt } (2*\text{pi}*n) * (n / \text{exp } 1)^n > a * s^{(2*n+1)}$ 
  using mult-strict-right-mono [OF a-less] mult-left-mono [OF s-le]
by (smt (verit, best) s-le ab(1) assms of-int-1 of-int-le-iff of-int-mult zero-less-power)
then have n: fact  $n > a * s^{(2*n+1)}$ 
  using fact-bounds(1) by (smt (verit, best)  $\langle 0 < n \rangle$  of-int-fact of-int-less-iff)
define F where  $F \equiv \lambda x. \sum_{i \leq 2*n.} (-1)^i * s^{(2*n-i)} * (\text{deriv } i) (hf\ n)\ x$ 
have Fder [derivative-intros]: (F has-real-derivative  $-s * F\ x + s^{(2*n+1)} * hf\ n\ x$ ) (at x) for x
proof –
  have *:  $\text{sum } f \{..n+n\} = \text{sum } f \{..<n+n\}$  if  $f\ (n+n) = 0$  for  $f::\text{nat} \Rightarrow \text{real}$ 
  by (smt (verit, best) lessThan-Suc-atMost sum.lessThan-Suc that)
  have [simp]:  $(\text{deriv } ((\text{deriv } (n+n)) (hf\ n))\ x) = 0$ 
  using hf-deriv-eq-0 [where  $k = \text{Suc}(n+n)$ ] by simp
  have §:  $(\sum_{k \leq n+n.} (-1)^k * ((\text{deriv } (\text{Suc } k)) (hf\ n))\ x * \text{of-int } s^{(n+n-k)})$ 

```

```

      + s * (∑ j=0..n+n. (-1) ^ j * ((deriv ~ j) (hf n) x * of-int s ^ (n+n
- j)))
    = s * (hf n x * of-int s ^ (n+n))
    using ⟨n>0⟩
    apply (subst sum-Suc-reindex)
    apply (simp add: algebra-simps atLeast0AtMost)
    apply (force simp add: * mult.left-commute [of of-int s] minus-nat.diff-Suc
sum-distrib-left
      simp flip: sum.distrib intro!: comm-monoid-add-class.sum.neutral
split: nat.split-asm)
    done
    show ?thesis
    unfolding F-def
    apply (rule derivative-eq-intros field-differentiable-derivI | simp)+
    using § by (simp add: algebra-simps atLeast0AtMost eval-nat-numeral)
qed

```

```

have F01-Ints: F 0 ∈ ℤ F 1 ∈ ℤ
  by (simp-all add: F-def hf-deriv-0 hf-deriv-1 Ints-sum)
define sF where sF ≡ λx. exp (of-int s * x) * F x
define sF' where sF' ≡ λx. of-int s ^ Suc(2*n) * (exp (of-int s * x) * hf n x)
have sF-der: (sF has-real-derivative sF' x) (at x) for x
  unfolding sF-def sF'-def
  by (rule refl derivative-eq-intros | force simp: algebra-simps)+
let ?N = b * integral {0..1} sF'
have sF'-integral: (sF' has-integral sF 1 - sF 0) {0..1}
  by (smt (verit) fundamental-theorem-of-calculus has-field-derivative-iff-has-vector-derivative
has-vector-derivative-at-within sF-der)
then have ?N = a * F 1 - b * F 0
  using ⟨b > 0⟩ by (simp add: integral-unique exp-s sF-def algebra-simps)
also have ... ∈ ℤ
  using hf-deriv-1 by (simp add: F01-Ints)
finally have N-Ints: ?N ∈ ℤ .
have sF' (1/2) > 0 and ge0: ∧x. x ∈ {0..1} ⇒ 0 ≤ sF' x
  using assms by (auto simp add: sF'-def hf-def)
moreover have continuous-on {0..1} sF'
  unfolding sF'-def hf-def by (intro continuous-intros) auto
ultimately have False if (sF' has-integral 0) {0..1}
  using has-integral-0-cbox-imp-0 [of 0 1 sF' 1/2] that by auto
then have integral {0..1} sF' > 0
  by (metis ge0 has-integral-nonneg integral-unique order-le-less sF'-integral)
then have 0 < ?N
  by (simp add: ⟨b > 0⟩)
have integral {0..1} sF' = of-int s ^ Suc(2*n) * integral {0..1} (λx. exp (s*x)
* hf n x)
  unfolding sF'-def by force
also have ... ≤ of-int s ^ Suc(2*n) * (exp s * (1 / fact n))
proof (rule mult-left-mono)
  have integral {0..1} (λx. exp (s*x) * hf n x) ≤ integral {0..1} (λx::real. exp s

```



```

* (1 / fact n)
  proof (intro mult-mono integral-le)
    show (λx. exp (s*x) * hf n x) integrable-on {0..1}
      using ‹0 < ?N› not-integrable-integral sF'-def by fastforce
  qed (use assms hf-nonneg hf-le-inverse-fact in auto)
  also have ... = exp s * (1 / fact n)
    by simp
  finally show integral {0..1} (λx. exp (s*x) * hf n x) ≤ exp s * (1 / fact n) .
  qed (use assms in auto)
  finally have ?N ≤ b * of-int s ^ Suc(2*n) * exp s * (1 / fact n)
    using ‹b > 0› by (simp add: sF'-def mult-ac divide-simps)
  also have ... < 1
    using n apply (simp add: field-simps exp-s)
    by (metis of-int-fact of-int-less-iff of-int-mult of-int-power)
  finally show False
    using ‹0 < ?N› Ints-cases N-Ints by force
qed

```

**theorem** *exp-irrational*:

```

fixes q::real assumes q ∈ ℚ q ≠ 0 shows exp q ∉ ℚ
proof
  assume q: exp q ∈ ℚ
  obtain s t where s ≠ 0 t > 0 q = of-int s / of-int t
    by (metis Rats-cases' assms div-0 of-int-0)
  then have (exp q) ^ (nat t) = exp s
    by (smt (verit, best) exp-divide-power-eq of-nat-nat zero-less-nat-eq)
  moreover have exp q ^ (nat t) ∈ ℚ
    by (simp add: q)
  ultimately show False
    by (smt (verit, del-insts) Rats-inverse ‹s ≠ 0› exp-minus exp-nat-irrational
of-int-of-nat)
qed

```

end

## References

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- [2] M. Eberl. Stirling’s formula. *Archive of Formal Proofs*, Sept. 2016. [https://isa-afp.org/entries/Stirling\\_Formula.html](https://isa-afp.org/entries/Stirling_Formula.html), Formal proof development.
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