

Hybrid Multi-Lane Spatial Logic

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Abstract

We present a semantic embedding of a spatio-temporal multi-modal logic, specifically defined to reason about motorway traffic, into Isabelle/HOL. The semantic model is an abstraction of a motorway, emphasising local spatial properties, and parameterised by the types of sensors deployed in the vehicles. We use the logic to define controller constraints to ensure safety, i.e., the absence of collisions on the motorway. After proving safety with a restrictive definition of sensors, we relax these assumptions and show how to amend the controller constraints to still guarantee safety.

Published in iFM 2017 [4].

Formal verification of autonomous vehicles on motorways is a challenging problem, due to the complex interactions between dynamical behaviours and controller choices of the vehicles. To overcome the complexities of proving safety properties, we proposed to separate the dynamical behaviour from the concrete changes in space [2]. To that end, we defined *Multi-Lane Spatial Logic* (MLSL), which was used to express guards and invariants of controller automata defining a protocol for safe lane-change manoeuvres. Under the assumption that all vehicles adhere to this protocol, we proved that collisions were avoided. Subsequently, we presented an extension of MLSL to reason about changes in space over time, a system of natural deduction, and formally proved a safety theorem [5, 3]. This proof was carried out manually and dependent on strong assumptions about the vehicles' sensors.

We define a semantic embedding of a further extension of MLSL, inspired by Hybrid Logic [1]. Subsequently, we show how the safety theorem can be proved within this embedding. Finally, we alter this formal embedding by relaxing the assumptions on the sensors. We show that the previously proven safety theorem does *not* ensure safety in this case, and how the controller constraints can be strengthened to guarantee safety.

Contents

1	Discrete Intervals based on Natural Numbers
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2

1.1	Basic properties of discrete intervals.	3
1.2	Algebraic properties of intersection and union.	6
2	Closed Real-valued Intervals	10
3	Cars	13
4	Traffic Snapshots	13
5	Views on Traffic	19
6	Restrict Claims and Reservations to a View	24
7	Move a View according to Difference between Traffic Snapshots	27
8	Sensors for Cars	27
9	Visible Length of Cars with Perfect Sensors	28
10	Basic HMLSL	31
10.1	Syntax of Basic HMLSL	31
10.2	Theorems about Basic HMLSL	34
11	Perfect Sensors	40
12	HMLSL for Perfect Sensors	41
13	Safety for Cars with Perfect Sensors	43
14	Regular Sensors	45
15	HMLSL for Regular Sensors	46
16	Safety for Cars with Regular Sensors	47

1 Discrete Intervals based on Natural Numbers

We define a type of intervals based on the natural numbers. To that end, we employ standard operators of Isabelle, but in addition prove some structural properties of the intervals. In particular, we show that this type constitutes a meet-semilattice with a bottom element and equality.

Furthermore, we show that this semilattice allows for a constrained join, i.e., the union of two intervals is defined, if either one of them is empty, or they are consecutive. Finally, we define the notion of *chopping* an interval into two consecutive subintervals.

```

theory NatInt
  imports Main
begin

```

A discrete interval is a set of consecutive natural numbers, or the empty set.

```

typedef nat-int = { S . (∃ (m::nat) n . {m..n }=S) }
  <proof>
setup-lifting type-definition-nat-int

```

1.1 Basic properties of discrete intervals.

```

locale nat-int
interpretation nat-int-class?: nat-int <proof>

```

```

context nat-int
begin

```

```

lemma un-consec-seq: (m::nat) ≤ n ∧ n+1 ≤ l ⟶ {m..n} ∪ {n+1..l} = {m..l}
  <proof>

```

```

lemma int-consec-seq: {(m::nat)..n} ∩ {n+1..l} = {}
  <proof>

```

```

lemma empty-type: {} ∈ { S . ∃ (m:: nat) n . {m..n}=S }
  <proof>

```

```

lemma inter-result: ∀ x ∈ { S . (∃ (m::nat) n . {m..n }=S) }.
  ∀ y ∈ { S . (∃ (m::nat) n . {m..n }=S) }.
  x ∩ y ∈ { S . (∃ (m::nat) n . {m..n }=S) }
  <proof>

```

```

lemma union-result: ∀ x ∈ { S . (∃ (m::nat) n . {m..n }=S) }.
  ∀ y ∈ { S . (∃ (m::nat) n . {m..n }=S) }.
  x ≠ {} ∧ y ≠ {} ∧ Max x + 1 = Min y
  ⟶ x ∪ y ∈ { S . (∃ (m::nat) n . {m..n }=S) }
  <proof>

```

```

lemma union-empty-result1: ∀ i ∈ { S . (∃ (m::nat) n . {m..n }=S) }.
  i ∪ {} ∈ { S . (∃ (m::nat) n . {m..n }=S) }
  <proof>

```

```

lemma union-empty-result2: ∀ i ∈ { S . (∃ (m::nat) n . {m..n }=S) }.
  {} ∪ i ∈ { S . (∃ (m::nat) n . {m..n }=S) }
  <proof>

```

```

lemma finite: ∀ i ∈ { S . (∃ (m::nat) n . {m..n }=S) } . (finite i)
  <proof>

```

```

lemma not-empty-means-seq:  $\forall i \in \{S . (\exists (m::nat) n . \{m..n\} = S) \} . i \neq \{\}$ 
 $\longrightarrow (\exists m n . m \leq n \wedge \{m..n\} = i)$ 
   $\langle proof \rangle$ 
end

```

The empty set is the bottom element of the type. The infimum/meet of the semilattice is set intersection. The order is given by the subset relation.

```

instantiation nat-int :: bot
begin
lift-definition bot-nat-int :: nat-int is Set.empty  $\langle proof \rangle$ 
instance  $\langle proof \rangle$ 
end

```

```

instantiation nat-int :: inf
begin
lift-definition inf-nat-int :: nat-int  $\Rightarrow$  nat-int  $\Rightarrow$  nat-int is Set.inter  $\langle proof \rangle$ 
instance
 $\langle proof \rangle$ 
end

```

```

instantiation nat-int :: order-bot
begin
lift-definition less-eq-nat-int :: nat-int  $\Rightarrow$  nat-int  $\Rightarrow$  bool is Set.subset-eq  $\langle proof \rangle$ 
lift-definition less-nat-int :: nat-int  $\Rightarrow$  nat-int  $\Rightarrow$  bool is Set.subset  $\langle proof \rangle$ 
instance
 $\langle proof \rangle$ 
end

```

```

instantiation nat-int :: semilattice-inf
begin
instance
 $\langle proof \rangle$ 
end

```

```

instantiation nat-int :: equal
begin
definition equal-nat-int :: nat-int  $\Rightarrow$  nat-int  $\Rightarrow$  bool
  where equal-nat-int i j  $\equiv i \leq j \wedge j \leq i$ 
instance
 $\langle proof \rangle$ 
end

```

```

context nat-int
begin
abbreviation subseq :: nat-int  $\Rightarrow$  nat-int  $\Rightarrow$  bool (infix  $\langle \sqsubseteq \rangle$  30)
  where i  $\sqsubseteq$  j  $\equiv i \leq j$ 
abbreviation empty :: nat-int ( $\langle \emptyset \rangle$ )

```

where $\emptyset \equiv \text{bot}$

notation inf (**infix** $\langle \sqcap \rangle$ 70)

The union of two intervals is only defined, if it is also a discrete interval.

definition $\text{union} :: \text{nat-int} \Rightarrow \text{nat-int} \Rightarrow \text{nat-int}$ (**infix** $\langle \sqcup \rangle$ 65)
where $i \sqcup j = \text{Abs-nat-int} (\text{Rep-nat-int } i \cup \text{Rep-nat-int } j)$

Non-empty intervals contain a minimal and maximal element. Two non-empty intervals i and j are consecutive, if the minimum of j is the successor of the maximum of i . Furthermore, the interval i can be chopped into the intervals j and k , if the union of j and k equals i , and if j and k are not-empty, they must be consecutive. Finally, we define the cardinality of discrete intervals by lifting the cardinality of sets.

definition $\text{maximum} :: \text{nat-int} \Rightarrow \text{nat}$
where $\text{maximum-def}: i \neq \emptyset \implies \text{maximum}(i) = \text{Max}(\text{Rep-nat-int } i)$

definition $\text{minimum} :: \text{nat-int} \Rightarrow \text{nat}$
where $\text{minimum-def}: i \neq \emptyset \implies \text{minimum}(i) = \text{Min}(\text{Rep-nat-int } i)$

definition $\text{consec} :: \text{nat-int} \Rightarrow \text{nat-int} \Rightarrow \text{bool}$
where $\text{consec } i \ j \equiv (i \neq \emptyset \wedge j \neq \emptyset \wedge (\text{maximum}(i)+1 = \text{minimum } j))$

definition $N\text{-Chop} :: \text{nat-int} \Rightarrow \text{nat-int} \Rightarrow \text{nat-int} \Rightarrow \text{bool}$ ($\langle N'\text{-Chop}'(-,-,-) \rangle$ 51)
where $n\text{chop-def} :$
 $N\text{-Chop}(i,j,k) \equiv (i = j \sqcup k \wedge (j = \emptyset \vee k = \emptyset \vee \text{consec } j \ k))$

lift-definition $\text{card}' :: \text{nat-int} \Rightarrow \text{nat}$ ($\langle |-| \rangle$ 70) **is** card $\langle \text{proof} \rangle$

For convenience, we also lift the membership relation and its negation to discrete intervals.

lift-definition $\text{el} :: \text{nat} \Rightarrow \text{nat-int} \Rightarrow \text{bool}$ (**infix** $\langle \in \rangle$ 50) **is** Set.member $\langle \text{proof} \rangle$

lift-definition $\text{not-in} :: \text{nat} \Rightarrow \text{nat-int} \Rightarrow \text{bool}$ (**infix** $\langle \notin \rangle$ 40) **is** Set.not-member $\langle \text{proof} \rangle$
end

lemmas $[\text{simp}] = \text{nat-int.el.rep-eq nat-int.not-in.rep-eq nat-int.card'.rep-eq}$

context nat-int
begin

lemma $\text{in-not-in-iff1} : n \in i \longleftrightarrow \neg n \notin i$ $\langle \text{proof} \rangle$
lemma $\text{in-not-in-iff2} : n \notin i \longleftrightarrow \neg n \in i$ $\langle \text{proof} \rangle$

lemma $\text{rep-non-empty-means-seq} : i \neq \emptyset$
 $\longrightarrow (\exists m \ n. m \leq n \wedge (\{m..n\} = (\text{Rep-nat-int } i)))$

$\langle proof \rangle$

lemma *non-empty-max*: $i \neq \emptyset \longrightarrow (\exists m . \text{maximum}(i) = m)$
 $\langle proof \rangle$

lemma *non-empty-min*: $i \neq \emptyset \longrightarrow (\exists m . \text{minimum}(i) = m)$
 $\langle proof \rangle$

lemma *minimum-in*: $i \neq \emptyset \longrightarrow \text{minimum } i \in i$
 $\langle proof \rangle$

lemma *maximum-in*: $i \neq \emptyset \longrightarrow \text{maximum } i \in i$
 $\langle proof \rangle$

lemma *non-empty-elem-in*: $i \neq \emptyset \longleftrightarrow (\exists n . n \in i)$
 $\langle proof \rangle$

lemma *leq-nat-non-empty*: $(m::nat) \leq n \longrightarrow \text{Abs-nat-int}\{m..n\} \neq \emptyset$
 $\langle proof \rangle$

lemma *leq-max-sup*: $(m::nat) \leq n \longrightarrow \text{Max } \{m..n\} = n$
 $\langle proof \rangle$

lemma *leq-min-inf*: $(m::nat) \leq n \longrightarrow \text{Min } \{m..n\} = m$
 $\langle proof \rangle$

lemma *leq-max-sup'*: $(m::nat) \leq n \longrightarrow \text{maximum}(\text{Abs-nat-int}\{m..n\}) = n$
 $\langle proof \rangle$

lemma *leq-min-inf'*: $(m::nat) \leq n \longrightarrow \text{minimum}(\text{Abs-nat-int}\{m..n\}) = m$
 $\langle proof \rangle$

lemma *in-refl*: $(n::nat) \in \text{Abs-nat-int } \{n\}$
 $\langle proof \rangle$

lemma *in-singleton*: $m \in \text{Abs-nat-int}\{n\} \longrightarrow m = n$
 $\langle proof \rangle$

1.2 Algebraic properties of intersection and union.

lemma *inter-empty1*: $(i::nat-int) \sqcap \emptyset = \emptyset$
 $\langle proof \rangle$

lemma *inter-empty2*: $\emptyset \sqcap (i::nat-int) = \emptyset$
 $\langle proof \rangle$

lemma *un-empty-absorb1*: $i \sqcup \emptyset = i$
 $\langle proof \rangle$

lemma *un-empty-absorb2*: $\emptyset \sqcup i = i$
 $\langle \text{proof} \rangle$

Most properties of the union of two intervals depends on them being consecutive, to ensure that their union exists.

lemma *consec-un*: $\text{consec } i \ j \wedge n \notin \text{Rep-nat-int}(i) \cup \text{Rep-nat-int } j$
 $\longrightarrow n \notin (i \sqcup j)$
 $\langle \text{proof} \rangle$

lemma *un-subset1*: $\text{consec } i \ j \longrightarrow i \sqsubseteq i \sqcup j$
 $\langle \text{proof} \rangle$

lemma *un-subset2*: $\text{consec } i \ j \longrightarrow j \sqsubseteq i \sqcup j$
 $\langle \text{proof} \rangle$

lemma *inter-distr1*: $\text{consec } j \ k \longrightarrow i \sqcap (j \sqcup k) = (i \sqcap j) \sqcup (i \sqcap k)$
 $\langle \text{proof} \rangle$

lemma *inter-distr2*: $\text{consec } j \ k \longrightarrow (j \sqcup k) \sqcap i = (j \sqcap i) \sqcup (k \sqcap i)$
 $\langle \text{proof} \rangle$

lemma *consec-un-not-elem1*: $\text{consec } i \ j \wedge n \notin i \sqcup j \longrightarrow n \notin i$
 $\langle \text{proof} \rangle$

lemma *consec-un-not-elem2*: $\text{consec } i \ j \wedge n \notin i \sqcup j \longrightarrow n \notin j$
 $\langle \text{proof} \rangle$

lemma *consec-un-element1*: $\text{consec } i \ j \wedge n \in i \longrightarrow n \in i \sqcup j$
 $\langle \text{proof} \rangle$

lemma *consec-un-element2*: $\text{consec } i \ j \wedge n \in j \longrightarrow n \in i \sqcup j$
 $\langle \text{proof} \rangle$

lemma *consec-lesser*: $\text{consec } i \ j \longrightarrow (\forall n \ m. (n \in i \wedge m \in j \longrightarrow n < m))$
 $\langle \text{proof} \rangle$

lemma *consec-in-exclusive1*: $\text{consec } i \ j \wedge n \in i \longrightarrow n \notin j$
 $\langle \text{proof} \rangle$

lemma *consec-in-exclusive2*: $\text{consec } i \ j \wedge n \in j \longrightarrow n \notin i$
 $\langle \text{proof} \rangle$

lemma *consec-un-max*: $\text{consec } i \ j \longrightarrow \text{maximum } j = \text{maximum } (i \sqcup j)$
 $\langle \text{proof} \rangle$

lemma *consec-un-min*: $\text{consec } i \ j \longrightarrow \text{minimum } i = \text{minimum } (i \sqcup j)$
 $\langle \text{proof} \rangle$

lemma *consec-un-defined*:

$consec\ i\ j \longrightarrow (Rep\text{-}nat\text{-}int\ (i \sqcup j) \in \{S . (\exists\ (m::nat)\ n . \{m..n\}=S)\})$
 $\langle proof \rangle$

lemma *consec-un-min-max:*

$consec\ i\ j \longrightarrow Rep\text{-}nat\text{-}int(i \sqcup j) = \{minimum\ i .. maximum\ j\}$
 $\langle proof \rangle$

lemma *consec-un-equality:*

$(consec\ i\ j \wedge k \neq \emptyset)$
 $\longrightarrow (minimum\ (i \sqcup j) = minimum\ (k) \wedge maximum\ (i \sqcup j) = maximum\ (k))$
 $\longrightarrow i \sqcup j = k$
 $\langle proof \rangle$

lemma *consec-trans-lesser:*

$consec\ i\ j \wedge consec\ j\ k \longrightarrow (\forall\ n\ m. (n \in i \wedge m \in k \longrightarrow n < m))$
 $\langle proof \rangle$

lemma *consec-inter-empty:* $consec\ i\ j \implies i \sqcap j = \emptyset$
 $\langle proof \rangle$

lemma *consec-intermediate1:* $consec\ j\ k \wedge consec\ i\ (j \sqcup k) \longrightarrow consec\ i\ j$
 $\langle proof \rangle$

lemma *consec-intermediate2:* $consec\ i\ j \wedge consec\ (i \sqcup j)\ k \longrightarrow consec\ j\ k$
 $\langle proof \rangle$

lemma *un-assoc:* $consec\ i\ j \wedge consec\ j\ k \longrightarrow (i \sqcup j) \sqcup k = i \sqcup (j \sqcup k)$
 $\langle proof \rangle$

lemma *consec-assoc1:* $consec\ j\ k \wedge consec\ i\ (j \sqcup k) \longrightarrow consec\ (i \sqcup j)\ k$
 $\langle proof \rangle$

lemma *consec-assoc2:* $consec\ i\ j \wedge consec\ (i \sqcup j)\ k \longrightarrow consec\ i\ (j \sqcup k)$
 $\langle proof \rangle$

lemma *consec-assoc-mult:*

$(i2=\emptyset \vee consec\ i1\ i2) \wedge (i3=\emptyset \vee consec\ i3\ i4) \wedge (consec\ (i1 \sqcup i2)\ (i3 \sqcup i4))$
 $\longrightarrow (i1 \sqcup i2) \sqcup (i3 \sqcup i4) = (i1 \sqcup (i2 \sqcup i3)) \sqcup i4$
 $\langle proof \rangle$

lemma *card-subset-le:* $i \sqsubseteq i' \longrightarrow |i| \leq |i'|$
 $\langle proof \rangle$

lemma *card-subset-less:* $(i::nat\text{-}int) < i' \longrightarrow |i| < |i'|$
 $\langle proof \rangle$

lemma *card-empty-zero:* $|\emptyset| = 0$
 $\langle proof \rangle$

lemma *card-non-empty-geq-one*: $i \neq \emptyset \longleftrightarrow |i| \geq 1$
 $\langle proof \rangle$

lemma *card-min-max*: $i \neq \emptyset \longrightarrow |i| = (\text{maximum } i - \text{minimum } i) + 1$
 $\langle proof \rangle$

lemma *card-un-add*: $\text{consec } i \ j \longrightarrow |i \sqcup j| = |i| + |j|$
 $\langle proof \rangle$

lemma *singleton*: $|i| = 1 \longrightarrow (\exists n. \text{Rep-nat-int } i = \{n\})$
 $\langle proof \rangle$

lemma *singleton2*: $(\exists n. \text{Rep-nat-int } i = \{n\}) \longrightarrow |i| = 1$
 $\langle proof \rangle$

lemma *card-seq*:
 $\forall i. |i| = x \longrightarrow (\text{Rep-nat-int } i = \{\} \vee (\exists n. \text{Rep-nat-int } i = \{n..n+(x-1)\}))$
 $\langle proof \rangle$

lemma *rep-single*: $\text{Rep-nat-int } (\text{Abs-nat-int } \{m..m\}) = \{m\}$
 $\langle proof \rangle$

lemma *chop-empty-right*: $\forall i. N\text{-Chop}(i, i, \emptyset)$
 $\langle proof \rangle$

lemma *chop-empty-left*: $\forall i. N\text{-Chop}(i, \emptyset, i)$
 $\langle proof \rangle$

lemma *chop-empty* : $N\text{-Chop}(\emptyset, \emptyset, \emptyset)$
 $\langle proof \rangle$

lemma *chop-always-possible*: $\forall i. \exists j \ k. N\text{-Chop}(i, j, k)$
 $\langle proof \rangle$

lemma *chop-add1*: $N\text{-Chop}(i, j, k) \longrightarrow |i| = |j| + |k|$
 $\langle proof \rangle$

lemma *chop-add2*: $|i| = x + y \longrightarrow (\exists j \ k. N\text{-Chop}(i, j, k) \wedge |j| = x \wedge |k| = y)$
 $\langle proof \rangle$

lemma *chop-single*: $(N\text{-Chop}(i, j, k) \wedge |i| = 1) \longrightarrow (|j| = 0 \vee |k| = 0)$
 $\langle proof \rangle$

lemma *chop-leq-max*: $N\text{-Chop}(i, j, k) \wedge \text{consec } j \ k \longrightarrow$
 $(\forall n. n \in \text{Rep-nat-int } i \wedge n \leq \text{maximum } j \longrightarrow n \in \text{Rep-nat-int } j)$
 $\langle proof \rangle$

lemma *chop-geq-min*: $N\text{-Chop}(i,j,k) \wedge \text{consec } j \ k \longrightarrow$
 $(\forall n . n \in \text{Rep-nat-int } i \wedge \text{minimum } k \leq n \longrightarrow n \in \text{Rep-nat-int } k)$
 $\langle \text{proof} \rangle$

lemma *chop-min*: $N\text{-Chop}(i,j,k) \wedge \text{consec } j \ k \longrightarrow \text{minimum } i = \text{minimum } j$
 $\langle \text{proof} \rangle$

lemma *chop-max*: $N\text{-Chop}(i,j,k) \wedge \text{consec } j \ k \longrightarrow \text{maximum } i = \text{maximum } k$
 $\langle \text{proof} \rangle$

lemma *chop-assoc1*:
 $N\text{-Chop}(i,i1,i2) \wedge N\text{-Chop}(i2,i3,i4)$
 $\longrightarrow (N\text{-Chop}(i, i1 \sqcup i3, i4) \wedge N\text{-Chop}(i1 \sqcup i3, i1, i3))$
 $\langle \text{proof} \rangle$

lemma *chop-assoc2*:
 $N\text{-Chop}(i,i1,i2) \wedge N\text{-Chop}(i1,i3,i4)$
 $\longrightarrow N\text{-Chop}(i, i3, i4 \sqcup i2) \wedge N\text{-Chop}(i4 \sqcup i2, i4, i2)$
 $\langle \text{proof} \rangle$

lemma *chop-subset1*: $N\text{-Chop}(i,j,k) \longrightarrow j \sqsubseteq i$
 $\langle \text{proof} \rangle$

lemma *chop-subset2*: $N\text{-Chop}(i,j,k) \longrightarrow k \sqsubseteq i$
 $\langle \text{proof} \rangle$

end
end

2 Closed Real-valued Intervals

We define a type for real-valued intervals. It consists of pairs of real numbers, where the first is lesser or equal to the second. Both endpoints are understood to be part of the interval, i.e., the intervals are closed. This also implies that we do not consider empty intervals.

We define a measure on these intervals as the difference between the left and right endpoint. In addition, we introduce a notion of shifting an interval by a real value x . Finally, an interval r can be chopped into s and t , if the left endpoint of r and s as well as the right endpoint of r and t coincides, and if the right endpoint of s is the left endpoint of t .

theory *RealInt*
imports *HOL.Real*
begin

typedef *real-int* = $\{r::(\text{real}*\text{real}) . \text{fst } r \leq \text{snd } r\}$
 $\langle \text{proof} \rangle$

setup-lifting *type-definition-real-int*

```

lift-definition left::real-int  $\Rightarrow$  real is fst  $\langle$ proof $\rangle$ 
lift-definition right::real-int  $\Rightarrow$  real is snd  $\langle$ proof $\rangle$ 

lemmas[simp] = left.rep-eq right.rep-eq

locale real-int
interpretation real-int-class?: real-int  $\langle$ proof $\rangle$ 

context real-int
begin

definition length :: real-int  $\Rightarrow$  real ( $\langle$ ||- $\rangle$  70)
  where  $\|r\| \equiv \text{right } r - \text{left } r$ 

definition shift::real-int  $\Rightarrow$  real  $\Rightarrow$  real-int ( $\langle$  shift -  $\rangle$ )
  where (shift r x) = Abs-real-int(left r + x, right r + x)

definition R-Chop :: real-int  $\Rightarrow$  real-int  $\Rightarrow$  real-int  $\Rightarrow$  bool ( $\langle$ R'-Chop'(-, -, -) $\rangle$  51)
  where rchop-def :
    R-Chop(r, s, t) == left r = left s  $\wedge$  right s = left t  $\wedge$  right r = right t

end

The intervals defined in this way allow for the definition of an order: the
subinterval relation.

instantiation real-int :: order
begin
definition less-eq-real-int r s  $\equiv$  (left r  $\geq$  left s)  $\wedge$  (right r  $\leq$  right s)
definition less-real-int r s  $\equiv$  (left r  $\geq$  left s)  $\wedge$  (right r  $\leq$  right s)
   $\wedge \neg((\text{left } s \geq \text{left } r) \wedge (\text{right } s \leq \text{right } r))$ 

instance
 $\langle$ proof $\rangle$ 
end

context real-int
begin

lemma left-leq-right: left r  $\leq$  right r
   $\langle$ proof $\rangle$ 

lemma length-ge-zero :  $\|r\| \geq 0$ 
   $\langle$ proof $\rangle$ 

lemma consec-add:
  left r = left s  $\wedge$  right r = right t  $\wedge$  right s = left t  $\implies \|r\| = \|s\| + \|t\|$ 
   $\langle$ proof $\rangle$ 

```

lemma *length-zero-iff-borders-eq*: $\|r\| = 0 \longleftrightarrow \text{left } r = \text{right } r$
 $\langle \text{proof} \rangle$

lemma *shift-left-eq-right*: $\text{left } (\text{shift } r \ x) \leq \text{right } (\text{shift } r \ x)$
 $\langle \text{proof} \rangle$

lemma *shift-keeps-length*: $\|r\| = \|\text{shift } r \ x\|$
 $\langle \text{proof} \rangle$

lemma *shift-zero*: $(\text{shift } r \ 0) = r$
 $\langle \text{proof} \rangle$

lemma *shift-additivity*: $(\text{shift } r \ (x+y)) = \text{shift } (\text{shift } r \ x) \ y$
 $\langle \text{proof} \rangle$

lemma *chop-always-possible*: $\forall r . \exists s \ t. R\text{-Chop}(r, s, t)$
 $\langle \text{proof} \rangle$

lemma *chop-singleton-right*: $\forall r . \exists s. R\text{-Chop}(r, r, s)$
 $\langle \text{proof} \rangle$

lemma *chop-singleton-left*: $\forall r . \exists s. R\text{-Chop}(r, s, r)$
 $\langle \text{proof} \rangle$

lemma *chop-add-length*: $R\text{-Chop}(r, s, t) \implies \|r\| = \|s\| + \|t\|$
 $\langle \text{proof} \rangle$

lemma *chop-add-length-ge-0*: $R\text{-Chop}(r, s, t) \wedge \|s\| > 0 \wedge \|t\| > 0 \longrightarrow \|r\| > 0$
 $\langle \text{proof} \rangle$

lemma *chop-dense* : $\|r\| > 0 \longrightarrow (\exists s \ t. R\text{-Chop}(r, s, t) \wedge \|s\| > 0 \wedge \|t\| > 0)$
 $\langle \text{proof} \rangle$

lemma *chop-assoc1*:
 $R\text{-Chop}(r, r1, r2) \wedge R\text{-Chop}(r2, r3, r4)$
 $\longrightarrow R\text{-Chop}(r, \text{Abs-real-int}(\text{left } r1, \text{right } r3), r4)$
 $\wedge R\text{-Chop}(\text{Abs-real-int}(\text{left } r1, \text{right } r3), r1, r3)$
 $\langle \text{proof} \rangle$

lemma *chop-assoc2*:
 $R\text{-Chop}(r, r1, r2) \wedge R\text{-Chop}(r1, r3, r4)$
 $\longrightarrow R\text{-Chop}(r, r3, \text{Abs-real-int}(\text{left } r4, \text{right } r2))$
 $\wedge R\text{-Chop}(\text{Abs-real-int}(\text{left } r4, \text{right } r2), r4, r2)$
 $\langle \text{proof} \rangle$

lemma *chop-leq1*: $R\text{-Chop}(r, s, t) \longrightarrow s \leq r$
 $\langle \text{proof} \rangle$

lemma *chop-leq2*: $R\text{-Chop}(r, s, t) \longrightarrow t \leq r$

<proof>

lemma *chop-empty1*: $R\text{-Chop}(r,s,t) \wedge \|s\| = 0 \longrightarrow r = t$
<proof>

lemma *chop-empty2*: $R\text{-Chop}(r,s,t) \wedge \|t\| = 0 \longrightarrow r = s$
<proof>

end

end

3 Cars

We define a type to refer to cars. For simplicity, we assume that (countably) infinite cars exist.

theory *Cars*
 imports *Main*
begin

The type of cars consists of the natural numbers. However, we do not define or prove any additional structure about it.

typedef *cars* = $\{n::nat. True\}$ *<proof>*

locale *cars*
begin

For the construction of possible counterexamples, it is beneficial to prove that at least two cars exist. Furthermore, we show that there indeed exist infinitely many cars.

lemma *at-least-two-cars-exists*: $\exists c\ d :: cars . c \neq d$
<proof>

lemma *infinite-cars*: *infinite* $\{c::cars . True\}$
<proof>

end
end

4 Traffic Snapshots

Traffic snapshots define the spatial and dynamical arrangement of cars on the whole of the motorway at a single point in time. A traffic snapshot consists of several functions assigning spatial properties and dynamical behaviour to each car. The functions are named as follows.

- pos: positions of cars
- res: reservations of cars
- clm: claims of cars
- dyn: current dynamic behaviour of cars
- physical_size: the real sizes of cars
- braking_distance: braking distance each car needs in emergency

```

theory Traffic
imports NatInt RealInt Cars
begin

```

```

type-synonym lanes = nat-int
type-synonym extension = real-int

```

Definition of the type of traffic snapshots. The constraints on the different functions are the *sanity conditions* of traffic snapshots.

```

typedef traffic =
  { ts :: (cars ⇒ real) * (cars ⇒ lanes) * (cars ⇒ lanes) * (cars ⇒ real ⇒ real) * (cars ⇒ real) * (cars ⇒ real).
    (∀ c. ((fst (snd ts))) c ⊓ ((fst (snd (snd ts)))) c = ∅) ∧
    (∀ c. |(fst (snd ts)) c| ≥ 1) ∧
    (∀ c. |(fst (snd ts)) c| ≤ 2) ∧
    (∀ c. |(fst (snd (snd ts)) c)| ≤ 1) ∧
    (∀ c. |(fst (snd ts)) c| + |(fst (snd (snd ts))) c| ≤ 2) ∧
    (∀ c. (fst (snd (snd (ts)))) c ≠ ∅ ⟶
      (∃ n. Rep-nat-int (fst (snd ts) c) ∪ Rep-nat-int (fst (snd (snd ts)) c)
        = {n, n+1})) ∧
    (∀ c. fst (snd (snd (snd (snd (ts)))))) c > 0) ∧
    (∀ c. snd (snd (snd (snd (snd (ts)))))) c > 0)
  }
  ⟨proof⟩

```

```

locale traffic
begin

```

```

notation nat-int.consec (↦consec↦)

```

For brevity, we define names for the different functions within a traffic snapshot.

```

definition pos::traffic ⇒ (cars ⇒ real)
where pos ts ≡ fst (Rep-traffic ts)

```

```

definition res::traffic ⇒ (cars ⇒ lanes)
where res ts ≡ fst (snd (Rep-traffic ts))

```

definition $clm :: traffic \Rightarrow (cars \Rightarrow lanes)$
where $clm\ ts \equiv fst\ (snd\ (snd\ (Rep-traffic\ ts)))$

definition $dyn :: traffic \Rightarrow (cars \Rightarrow (real \Rightarrow real))$
where $dyn\ ts \equiv fst\ (snd\ (snd\ (snd\ (Rep-traffic\ ts))))$

definition $physical-size :: traffic \Rightarrow (cars \Rightarrow real)$
where $physical-size\ ts \equiv fst\ (snd\ (snd\ (snd\ (Rep-traffic\ ts))))$

definition $braking-distance :: traffic \Rightarrow (cars \Rightarrow real)$
where $braking-distance\ ts \equiv snd\ (snd\ (snd\ (snd\ (Rep-traffic\ ts))))$

It is helpful to be able to refer to the sanity conditions of a traffic snapshot via lemmas, hence we prove that the sanity conditions hold for each traffic snapshot.

lemma *disjoint*: $(res\ ts\ c) \sqcap (clm\ ts\ c) = \emptyset$
 $\langle proof \rangle$

lemma *atLeastOneRes*: $1 \leq |res\ ts\ c|$
 $\langle proof \rangle$

lemma *atMostTwoRes*: $|res\ ts\ c| \leq 2$
 $\langle proof \rangle$

lemma *atMostOneCln*: $|clm\ ts\ c| \leq 1$
 $\langle proof \rangle$

lemma *atMostTwoLanes*: $|res\ ts\ c| + |clm\ ts\ c| \leq 2$
 $\langle proof \rangle$

lemma *consecutiveRes*: $|res\ ts\ c| = 2 \longrightarrow (\exists n. Rep-nat-int\ (res\ ts\ c) = \{n, n+1\})$
 $\langle proof \rangle$

lemma *clmNextRes* :
 $(clm\ ts\ c) \neq \emptyset \longrightarrow (\exists n. Rep-nat-int(res\ ts\ c) \cup Rep-nat-int(clm\ ts\ c) = \{n, n+1\})$
 $\langle proof \rangle$

lemma *psGeZero*: $\forall c. (physical-size\ ts\ c > 0)$
 $\langle proof \rangle$

lemma *sdGeZero*: $\forall c. (braking-distance\ ts\ c > 0)$
 $\langle proof \rangle$

While not a sanity condition directly, the following lemma helps to establish general properties of HMLSL later on. It is a consequence of *clmNextRes*.

lemma *clm-consec-res*:
 $(clm\ ts)\ c \neq \emptyset \longrightarrow consec\ (clm\ ts\ c)\ (res\ ts\ c) \vee consec\ (res\ ts\ c)\ (clm\ ts\ c)$
 $\langle proof \rangle$

We define several possible transitions between traffic snapshots. Cars may create or withdraw claims and reservations, as long as the sanity conditions of the traffic snapshots are fulfilled.

In particular, a car can only create a claim, if it possesses only a reservation on a single lane, and does not already possess a claim. Withdrawing a claim can be done in any situation. It only has an effect, if the car possesses a claim. Similarly, the transition for a car to create a reservation is always possible, but only changes the spatial situation on the road, if the car already has a claim. Finally, a car may withdraw its reservation to a single lane, if its current reservation consists of two lanes.

All of these transitions concern the spatial properties of a single car at a time, i.e., for several cars to change their properties, several transitions have to be taken.

definition *create-claim* ::

$$\begin{aligned}
& \text{traffic} \Rightarrow \text{cars} \Rightarrow \text{nat} \Rightarrow \text{traffic} \Rightarrow \text{bool} \quad (\langle - \rangle - c'(-, -') \rightarrow - \rangle 27) \\
\text{where} \quad (ts - c(c, n) \rightarrow ts') & == (pos\ ts') = (pos\ ts) \\
& \wedge (res\ ts') = (res\ ts) \\
& \wedge (dyn\ ts') = (dyn\ ts) \\
& \wedge (physical\text{-}size\ ts') = (physical\text{-}size\ ts) \\
& \wedge (braking\text{-}distance\ ts') = (braking\text{-}distance\ ts) \\
& \wedge |clm\ ts\ c| = 0 \\
& \wedge |res\ ts\ c| = 1 \\
& \wedge ((n+1) \in res\ ts\ c \vee (n-1 \in res\ ts\ c)) \\
& \wedge (clm\ ts') = (clm\ ts)(c := Abs\text{-}nat\text{-}int\ \{n\})
\end{aligned}$$

definition *withdraw-claim* ::

$$\begin{aligned}
& \text{traffic} \Rightarrow \text{cars} \Rightarrow \text{traffic} \Rightarrow \text{bool} \quad (\langle - \rangle - wdc'(-, -') \rightarrow - \rangle 27) \\
\text{where} \quad (ts - wdc(c) \rightarrow ts') & == (pos\ ts') = (pos\ ts) \\
& \wedge (res\ ts') = (res\ ts) \\
& \wedge (dyn\ ts') = (dyn\ ts) \\
& \wedge (physical\text{-}size\ ts') = (physical\text{-}size\ ts) \\
& \wedge (braking\text{-}distance\ ts') = (braking\text{-}distance\ ts) \\
& \wedge (clm\ ts') = (clm\ ts)(c := \emptyset)
\end{aligned}$$

definition *create-reservation* ::

$$\begin{aligned}
& \text{traffic} \Rightarrow \text{cars} \Rightarrow \text{traffic} \Rightarrow \text{bool} \quad (\langle - \rangle - r'(-, -') \rightarrow - \rangle 27) \\
\text{where} \quad (ts - r(c) \rightarrow ts') & == (pos\ ts') = (pos\ ts) \\
& \wedge (res\ ts') = (res\ ts)(c := (res\ ts\ c) \sqcup (clm\ ts\ c)) \\
& \wedge (dyn\ ts') = (dyn\ ts) \\
& \wedge (clm\ ts') = (clm\ ts)(c := \emptyset) \\
& \wedge (physical\text{-}size\ ts') = (physical\text{-}size\ ts) \\
& \wedge (braking\text{-}distance\ ts') = (braking\text{-}distance\ ts)
\end{aligned}$$

definition *withdraw-reservation* ::

$$\begin{aligned}
& \text{traffic} \Rightarrow \text{cars} \Rightarrow \text{nat} \Rightarrow \text{traffic} \Rightarrow \text{bool} \quad (\langle - \rangle - wdr'(-, -') \rightarrow - \rangle 27) \\
\text{where} \quad (ts - wdr(c, n) \rightarrow ts') & == (pos\ ts') = (pos\ ts)
\end{aligned}$$

$$\begin{aligned}
& \wedge (res\ ts') = (res\ ts)(c := Abs\text{-}nat\text{-}int\{n\}) \\
& \wedge (dyn\ ts') = (dyn\ ts) \\
& \wedge (clm\ ts') = (clm\ ts) \\
& \wedge (physical\text{-}size\ ts') = (physical\text{-}size\ ts) \\
& \wedge (braking\text{-}distance\ ts') = (braking\text{-}distance\ ts) \\
& \wedge n \in (res\ ts\ c) \\
& \wedge |res\ ts\ c| = 2
\end{aligned}$$

The following two transitions concern the dynamical behaviour of the cars. Similar to the spatial properties, a car may change its dynamics, by setting it to a new function f from real to real. Observe that this function is indeed arbitrary and does not constrain the possible behaviour in any way. However, this transition allows a car to change the function determining their braking distance (in fact, all cars are allowed to change this function, if a car changes sets a new dynamical function). That is, our model describes an over-approximation of a concrete situation, where the braking distance is determined by the dynamics.

The final transition describes the passing of x time units. That is, all cars update their position according to their current dynamical behaviour. Observe that this transition requires that the dynamics of each car is at least 0, for each time point between 0 and x . Hence, this condition denotes that all cars drive into the same direction. If the current dynamics of a car violated this constraint, it would have to reset its dynamics, until time may pass again.

definition *change-dyn::*

$$\begin{aligned}
& traffic \Rightarrow cars \Rightarrow (real \Rightarrow real) \Rightarrow traffic \Rightarrow bool\ (\langle - - dyn'(-, -) \rightarrow - \rangle 27) \\
\textbf{where } & (ts - dyn(c, f) \rightarrow ts') == (pos\ ts' = pos\ ts) \\
& \wedge (res\ ts' = res\ ts) \\
& \wedge (clm\ ts' = clm\ ts) \\
& \wedge (dyn\ ts' = (dyn\ ts)(c := f)) \\
& \wedge (physical\text{-}size\ ts') = (physical\text{-}size\ ts)
\end{aligned}$$

definition *drive::*

$$\begin{aligned}
& traffic \Rightarrow real \Rightarrow traffic \Rightarrow bool\ (\langle - - - \rightarrow - \rangle 27) \\
\textbf{where } & (ts - x \rightarrow ts') == (\forall c. (pos\ ts'\ c = (pos\ ts\ c) + (dyn\ ts\ c\ x))) \\
& \wedge (\forall c\ y. 0 \leq y \wedge y \leq x \longrightarrow dyn\ ts\ c\ y \geq 0) \\
& \wedge (res\ ts' = res\ ts) \\
& \wedge (clm\ ts' = clm\ ts) \\
& \wedge (dyn\ ts' = dyn\ ts) \\
& \wedge (physical\text{-}size\ ts') = (physical\text{-}size\ ts) \\
& \wedge (braking\text{-}distance\ ts') = (braking\text{-}distance\ ts)
\end{aligned}$$

We bundle the dynamical transitions into *evolutions*, since we will only reason about combinations of the dynamical behaviour. This fits to the level of abstraction by hiding the dynamics completely inside of the model.

inductive *evolve::* $traffic \Rightarrow traffic \Rightarrow bool\ (\langle - \rightsquigarrow - \rangle)$

where $refl : ts \rightsquigarrow ts \mid$
 $change: \exists c. \exists f. (ts - dyn(c, f) \rightarrow ts') \implies ts' \rightsquigarrow ts'' \implies ts \rightsquigarrow ts'' \mid$
 $drive: \exists x. x \geq 0 \wedge (ts - x \rightarrow ts') \implies ts' \rightsquigarrow ts'' \implies ts \rightsquigarrow ts''$

lemma $evolve-trans: (ts0 \rightsquigarrow ts1) \implies (ts1 \rightsquigarrow ts2) \implies (ts0 \rightsquigarrow ts2)$
 $\langle proof \rangle$

For general transition sequences, we introduce *abstract transitions*. A traffic snapshot ts' is reachable from ts via an abstract transition, if there is an arbitrary sequence of transitions from ts to ts' .

inductive $abstract::traffic \Rightarrow traffic \Rightarrow bool \ (\vdash \Rightarrow \rightarrow)$ **for** ts
where $refl: (ts \Rightarrow ts) \mid$

$evolve: ts \Rightarrow ts' \implies ts' \rightsquigarrow ts'' \implies ts \Rightarrow ts'' \mid$
 $cr-clm: ts \Rightarrow ts' \implies \exists c. \exists n. (ts' - c(c, n) \rightarrow ts'') \implies ts \Rightarrow ts'' \mid$
 $wd-clm: ts \Rightarrow ts' \implies \exists c. (ts' - wdc(c) \rightarrow ts'') \implies ts \Rightarrow ts'' \mid$
 $cr-res: ts \Rightarrow ts' \implies \exists c. (ts' - r(c) \rightarrow ts'') \implies ts \Rightarrow ts'' \mid$
 $wd-res: ts \Rightarrow ts' \implies \exists c. \exists n. (ts' - wdr(c, n) \rightarrow ts'') \implies ts \Rightarrow ts''$

lemma $abs-trans: (ts1 \Rightarrow ts2) \implies (ts0 \Rightarrow ts1) \implies (ts0 \Rightarrow ts2)$
 $\langle proof \rangle$

Most properties of the transitions are straightforward. However, to show that the transition to create a reservation is always possible, we need to explicitly construct the resulting traffic snapshot. Due to the size of such a snapshot, the proof is lengthy.

lemma $create-res-subseteq1: (ts - r(c) \rightarrow ts') \longrightarrow res\ ts\ c \sqsubseteq res\ ts'\ c$
 $\langle proof \rangle$

lemma $create-res-subseteq2: (ts - r(c) \rightarrow ts') \longrightarrow clm\ ts\ c \sqsubseteq res\ ts'\ c$
 $\langle proof \rangle$

lemma $create-res-subseteq1-neq: (ts - r(d) \rightarrow ts') \wedge d \neq c \longrightarrow res\ ts\ c = res\ ts'\ c$
 $\langle proof \rangle$

lemma $create-res-subseteq2-neq: (ts - r(d) \rightarrow ts') \wedge d \neq c \longrightarrow clm\ ts\ c = clm\ ts'\ c$
 $\langle proof \rangle$

lemma $always-create-res: \forall ts. \exists ts'. (ts - r(c) \rightarrow ts')$
 $\langle proof \rangle$

lemma $create-clm-eq-res: (ts - c(d, n) \rightarrow ts') \longrightarrow res\ ts\ c = res\ ts'\ c$
 $\langle proof \rangle$

lemma $withdraw-clm-eq-res: (ts - wdc(d) \rightarrow ts') \longrightarrow res\ ts\ c = res\ ts'\ c$
 $\langle proof \rangle$

lemma *withdraw-res-subseteq*: $(ts - wdr(d, n) \rightarrow ts') \longrightarrow res\ ts'\ c \sqsubseteq res\ ts\ c$
 $\langle proof \rangle$

end
end

5 Views on Traffic

In this section, we define a notion of locality for each car. These local parts of a road are called *views* and define the part of the model currently under consideration by a car. In particular, a view consists of

- the *extension*, a real-valued interval denoting the distance perceived,
- the *lanes*, a discrete interval, denoting which lanes are perceived,
- the *owner*, the car associated with this view.

theory *Views*
imports *NatInt RealInt Cars*
begin

type-synonym *lanes* = *nat-int*
type-synonym *extension* = *real-int*

record *view* =
ext::*extension*
lan::*lanes*
own::*cars*

The orders on discrete and continuous intervals induce an order on views. For two views v and v' with $v \leq v'$, we call v a *subview* of v' .

instantiation *view-ext*:: (*order*) *order*
begin

definition *less-eq-view-ext* ($V :: 'a\ view-ext$) ($V' :: 'a\ view-ext$) \equiv
 $(ext\ V \leq ext\ V') \wedge (lan\ V \sqsubseteq lan\ V') \wedge own\ V = own\ V'$
 $\wedge more\ V \leq more\ V'$

definition *less-view-ext* ($V :: 'a\ view-ext$) ($V' :: 'a\ view-ext$) \equiv
 $(ext\ V \leq ext\ V') \wedge (lan\ V \sqsubseteq lan\ V') \wedge own\ V' = own\ V$
 $\wedge more\ V \leq more\ V' \wedge$
 $\neg((ext\ V' \leq ext\ V) \wedge (lan\ V' \sqsubseteq lan\ V) \wedge own\ V' = own\ V$
 $\wedge more\ V' \leq more\ V)$

instance
 $\langle proof \rangle$
end

locale *view*

begin

notation *nat-int.maximum* ($\langle \text{maximum} \rangle$)

notation *nat-int.minimum* ($\langle \text{minimum} \rangle$)

notation *nat-int.consec* ($\langle \text{consec} \rangle$)

We lift the chopping relations from discrete and continuous intervals to views, and introduce new notation for these relations.

definition $hchop :: view \Rightarrow view \Rightarrow view \Rightarrow bool$ ($\langle \text{--}=\|\text{--} \rangle$)

where $(v=u\|w) == real\text{-}int.R\text{-}Chop(ext\ v)(ext\ u)(ext\ w) \wedge$
 $lan\ v=lan\ u \wedge$
 $lan\ v=lan\ w \wedge$
 $own\ v = own\ u \wedge$
 $own\ v = own\ w \wedge$
 $more\ v = more\ w \wedge$
 $more\ v = more\ u$

definition $vchop :: view \Rightarrow view \Rightarrow view \Rightarrow bool$ ($\langle \text{--}---\text{--} \rangle$)

where $(v=u---w) == nat\text{-}int.N\text{-}Chop(lan\ v)(lan\ u)(lan\ w) \wedge$
 $ext\ v = ext\ u \wedge$
 $ext\ v = ext\ w \wedge$
 $own\ v = own\ u \wedge$
 $own\ v = own\ w \wedge$
 $more\ v = more\ w \wedge$
 $more\ v = more\ u$

We can also switch the perspective of a view to the car *c*. That is, we substitute *c* for the original owner of the view.

definition $switch :: view \Rightarrow cars \Rightarrow view \Rightarrow bool$ ($\langle \text{--} = - > \text{--} \rangle$)

where $(v=c>w) == ext\ v = ext\ w \wedge$
 $lan\ v = lan\ w \wedge$
 $own\ w = c \wedge$
 $more\ v = more\ w$

Most of the lemmas in this theory are direct transfers of the corresponding lemmas on discrete and continuous intervals, which implies rather simple proofs. The only exception is the connection between subviews and the chopping operations. This proof is rather lengthy, since we need to distinguish several cases, and within each case, we need to explicitly construct six different views for the chopping relations.

lemma *h-chop-middle1*: $(v=u\|w) \longrightarrow left\ (ext\ v) \leq right\ (ext\ u)$
 $\langle \text{proof} \rangle$

lemma *h-chop-middle2*: $(v=u\|w) \longrightarrow right\ (ext\ v) \geq left\ (ext\ w)$
 $\langle \text{proof} \rangle$

lemma *horizontal-chop1*: $\exists u w. (v=u||w)$
 $\langle proof \rangle$

lemma *horizontal-chop-empty-right*: $\forall v. \exists u. (v=v||u)$
 $\langle proof \rangle$

lemma *horizontal-chop-empty-left*: $\forall v. \exists u. (v=u||v)$
 $\langle proof \rangle$

lemma *horizontal-chop-non-empty*:
 $\|ext\ v\| > 0 \longrightarrow (\exists u w. (v=u||w) \wedge \|ext\ u\| > 0 \wedge \|ext\ w\| > 0)$
 $\langle proof \rangle$

lemma *horizontal-chop-split-add*:
 $x \geq 0 \wedge y \geq 0 \longrightarrow \|ext\ v\| = x+y \longrightarrow (\exists u w. (v=u||w) \wedge \|ext\ u\| = x \wedge \|ext\ w\| = y)$
 $\langle proof \rangle$

lemma *horizontal-chop-assoc1*:
 $(v=v1||v2) \wedge (v2=v3||v4) \longrightarrow (\exists v'. (v=v'||v4) \wedge (v'=v1||v3))$
 $\langle proof \rangle$

lemma *horizontal-chop-assoc2*:
 $(v=v1||v2) \wedge (v1=v3||v4) \longrightarrow (\exists v'. (v=v3||v') \wedge (v'=v4||v2))$
 $\langle proof \rangle$

lemma *horizontal-chop-width-stable*: $(v=u||w) \longrightarrow |lan\ v| = |lan\ u| \wedge |lan\ v| = |lan\ w|$
 $\langle proof \rangle$

lemma *horizontal-chop-own-trans*: $(v=u||w) \longrightarrow own\ u = own\ w$
 $\langle proof \rangle$

lemma *vertical-chop1*: $\forall v. \exists u w. (v=u--w)$
 $\langle proof \rangle$

lemma *vertical-chop-empty-down*: $\forall v. \exists u. (v=v--u)$
 $\langle proof \rangle$

lemma *vertical-chop-empty-up*: $\forall v. \exists u. (v=u--v)$
 $\langle proof \rangle$

lemma *vertical-chop-assoc1*:
 $(v=v1--v2) \wedge (v2=v3--v4) \longrightarrow (\exists v'. (v=v'--v4) \wedge (v'=v1--v3))$

$\langle proof \rangle$

lemma *vertical-chop-assoc2*:

$(v=v1--v2) \wedge (v1=v3--v4) \longrightarrow (\exists v'. (v=v3--v') \wedge (v'=v4--v2))$
 $\langle proof \rangle$

lemma *vertical-chop-singleton*:

$(v=u--w) \wedge |lan\ v| = 1 \longrightarrow (|lan\ u| = 0 \vee |lan\ w| = 0)$
 $\langle proof \rangle$

lemma *vertical-chop-add1*: $(v=u--w) \longrightarrow |lan\ v| = |lan\ u| + |lan\ w|$
 $\langle proof \rangle$

lemma *vertical-chop-add2*:

$|lan\ v| = x+y \longrightarrow (\exists\ u\ w. (v=u--w) \wedge |lan\ u| = x \wedge |lan\ w| = y)$
 $\langle proof \rangle$

lemma *vertical-chop-length-stable*:

$(v=u--w) \longrightarrow \|ext\ v\| = \|ext\ u\| \wedge \|ext\ v\| = \|ext\ w\|$
 $\langle proof \rangle$

lemma *vertical-chop-own-trans*: $(v=u--w) \longrightarrow own\ u = own\ w$
 $\langle proof \rangle$

lemma *vertical-chop-width-mon*:

$(v=v1--v2) \wedge (v2=v3--v4) \wedge |lan\ v3| = x \longrightarrow |lan\ v| \geq x$
 $\langle proof \rangle$

lemma *horizontal-chop-leq1*: $(v=u\|w) \longrightarrow u \leq v$
 $\langle proof \rangle$

lemma *horizontal-chop-leq2*: $(v=u\|w) \longrightarrow w \leq v$
 $\langle proof \rangle$

lemma *vertical-chop-leq1*: $(v=u--w) \longrightarrow u \leq v$
 $\langle proof \rangle$

lemma *vertical-chop-leq2*: $(v=u--w) \longrightarrow w \leq v$
 $\langle proof \rangle$

lemma *somewhere-leq*:

$v \leq v' \iff (\exists\ v1\ v2\ v3\ v4\ v5\ v6. (v'=v1\|v2) \wedge (v1=v3--v4) \wedge (v2=v5--v6) \wedge (v3=v4--v5) \wedge (v4=v6--v1))$
 $\langle proof \rangle$

The switch relation is compatible with the chopping relations, in the following sense. If we can chop a view v into two subviews u and w , and we can

reach v' via the switch relation, then there also exist two subviews u', w' of v' , such that u' is reachable from u (and respectively for w', w).

lemma *switch-unique*: $(v = c > u) \wedge (v = c > w) \longrightarrow u = w$
 $\langle proof \rangle$

lemma *switch-exists*: $\exists c. u. (v = c > u)$
 $\langle proof \rangle$

lemma *switch-always-exists*: $\forall c. \exists u. (v = c > u)$
 $\langle proof \rangle$

lemma *switch-origin*: $\exists u. (u = (own\ v) > v)$
 $\langle proof \rangle$

lemma *switch-refl*: $(v = (own\ v) > v)$
 $\langle proof \rangle$

lemma *switch-symm*: $(v = c > u) \longrightarrow (u = (own\ v) > v)$
 $\langle proof \rangle$

lemma *switch-trans*: $(v = c > u) \wedge (u = d > w) \longrightarrow (v = d > w)$
 $\langle proof \rangle$

lemma *switch-triangle*: $(v = c > u) \wedge (v = d > w) \longrightarrow (u = d > w)$
 $\langle proof \rangle$

lemma *switch-hchop1*:
 $(v = v1 \parallel v2) \wedge (v = c > v') \longrightarrow$
 $(\exists v1' v2'. (v1 = c > v1') \wedge (v2 = c > v2') \wedge (v' = v1' \parallel v2'))$
 $\langle proof \rangle$

lemma *switch-hchop2*:
 $(v' = v1' \parallel v2') \wedge (v = c > v') \longrightarrow$
 $(\exists v1 v2. (v1 = c > v1') \wedge (v2 = c > v2') \wedge (v = v1 \parallel v2))$
 $\langle proof \rangle$

lemma *switch-vchop1*:
 $(v = v1 \dashv\vdash v2) \wedge (v = c > v') \longrightarrow$
 $(\exists v1' v2'. (v1 = c > v1') \wedge (v2 = c > v2') \wedge (v' = v1' \dashv\vdash v2'))$
 $\langle proof \rangle$

lemma *switch-vchop2*:
 $(v' = v1' \dashv\vdash v2') \wedge (v = c > v') \longrightarrow$
 $(\exists v1 v2. (v1 = c > v1') \wedge (v2 = c > v2') \wedge (v = v1 \dashv\vdash v2))$
 $\langle proof \rangle$

lemma *switch-leq*: $u' \leq u \wedge (v = c > u) \longrightarrow (\exists v'. (v' = c > u') \wedge v' \leq v)$
 $\langle proof \rangle$

end

end

6 Restrict Claims and Reservations to a View

To model that a view restricts the number of lanes a car may perceive, we define a function *restrict* taking a view v , a function f from cars to lanes and a car c , and returning the intersection between $f(c)$ and the set of lanes of v . This function will in the following only be applied to the functions yielding reservations and claims from traffic snapshots.

The lemmas of this section describe the connection between *restrict* and the different operations on traffic snapshots and views (e.g., the transition relations or the fact that reservations and claims are consecutive).

theory *Restriction*

imports *Traffic Views*

begin

locale *restriction* = *view+traffic*

begin

definition *restrict* :: $view \Rightarrow (cars \Rightarrow lanes) \Rightarrow cars \Rightarrow lanes$

where $restrict\ v\ f\ c == (f\ c) \sqcap lan\ v$

lemma *restrict-def'*: $restrict\ v\ f\ c = lan\ v \sqcap f\ c$

<proof>

lemma *restrict-subseteq*: $restrict\ v\ f\ c \sqsubseteq f\ c$

<proof>

lemma *restrict-clm* : $restrict\ v\ (clm\ ts)\ c \sqsubseteq clm\ ts\ c$

<proof>

lemma *restrict-res*: $restrict\ v\ (res\ ts)\ c \sqsubseteq res\ ts\ c$

<proof>

lemma *restrict-view*: $restrict\ v\ f\ c \sqsubseteq lan\ v$

<proof>

lemma *restriction-stable*: $(v=u \parallel w) \longrightarrow restrict\ u\ f\ c = restrict\ w\ f\ c$

<proof>

lemma *restriction-stable1*: $(v=u \parallel w) \longrightarrow restrict\ v\ f\ c = restrict\ u\ f\ c$

<proof>

lemma *restriction-stable2*: $(v=u \parallel w) \longrightarrow restrict\ v\ f\ c = restrict\ w\ f\ c$

<proof>

lemma *restriction-un*:

$$(v=u--w) \longrightarrow \text{restrict } v \text{ } f \text{ } c = (\text{restrict } u \text{ } f \text{ } c \sqcup \text{restrict } w \text{ } f \text{ } c)$$

$\langle \text{proof} \rangle$

lemma *restriction-mon1*: $(v=u--w) \longrightarrow \text{restrict } u \text{ } f \text{ } c \sqsubseteq \text{restrict } v \text{ } f \text{ } c$

$\langle \text{proof} \rangle$

lemma *restriction-mon2*: $(v=u--w) \longrightarrow \text{restrict } w \text{ } f \text{ } c \sqsubseteq \text{restrict } v \text{ } f \text{ } c$

$\langle \text{proof} \rangle$

lemma *restriction-disj*: $(v=u--w) \longrightarrow (\text{restrict } u \text{ } f \text{ } c) \sqcap (\text{restrict } w \text{ } f \text{ } c) = \emptyset$

$\langle \text{proof} \rangle$

lemma *vertical-chop-restriction-res-consec-or-empty*:

$$(v=v1--v2) \wedge \text{restrict } v1 \text{ } (res \text{ } ts) \text{ } c \neq \emptyset \wedge \text{consec } ((lan \text{ } v1)) ((lan \text{ } v2)) \wedge \\ \neg \text{consec } (\text{restrict } v1 \text{ } (res \text{ } ts) \text{ } c) (\text{restrict } v2 \text{ } (res \text{ } ts) \text{ } c) \\ \longrightarrow \text{restrict } v2 \text{ } (res \text{ } ts) \text{ } c = \emptyset$$

$\langle \text{proof} \rangle$

lemma *restriction-consec-res*: $(v=u--w)$

$$\longrightarrow \text{restrict } u \text{ } (res \text{ } ts) \text{ } c = \emptyset \vee \text{restrict } w \text{ } (res \text{ } ts) \text{ } c = \emptyset \\ \vee \text{consec } (\text{restrict } u \text{ } (res \text{ } ts) \text{ } c) (\text{restrict } w \text{ } (res \text{ } ts) \text{ } c)$$

$\langle \text{proof} \rangle$

lemma *restriction-clm-res-disjoint*:

$$(\text{restrict } v \text{ } (res \text{ } ts) \text{ } c) \sqcap (\text{restrict } v \text{ } (clm \text{ } ts) \text{ } c) = \emptyset$$

$\langle \text{proof} \rangle$

lemma *el-in-restriction-clm-singleton*:

$$n \in \text{restrict } v \text{ } (clm \text{ } ts) \text{ } c \longrightarrow (clm \text{ } ts) \text{ } c = \text{Abs-nat-int}(\{n\})$$

$\langle \text{proof} \rangle$

lemma *restriction-clm-v2-non-empty-v1-empty*:

$$(v=u--w) \wedge \text{restrict } w \text{ } (clm \text{ } ts) \text{ } c \neq \emptyset \wedge \\ \text{consec } ((lan \text{ } u)) ((lan \text{ } w)) \longrightarrow \text{restrict } u \text{ } (clm \text{ } ts) \text{ } c = \emptyset$$

$\langle \text{proof} \rangle$

lemma *restriction-consec-clm*:

$$(v=u--w) \wedge \text{consec } (lan \text{ } u) (lan \text{ } w) \\ \longrightarrow \text{restrict } u \text{ } (clm \text{ } ts) \text{ } c = \emptyset \vee \text{restrict } w \text{ } (clm \text{ } ts) \text{ } c = \emptyset$$

$\langle \text{proof} \rangle$

lemma *restriction-add-res*:

$$(v=u--w) \\ \longrightarrow |\text{restrict } v \text{ } (res \text{ } ts) \text{ } c| = |\text{restrict } u \text{ } (res \text{ } ts) \text{ } c| + |\text{restrict } w \text{ } (res \text{ } ts) \text{ } c|$$

$\langle \text{proof} \rangle$

lemma *restriction-eq-view-card*: $\text{restrict } v \ f \ c = \text{lan } v \longrightarrow |\text{restrict } v \ f \ c| = |\text{lan } v|$
 $\langle \text{proof} \rangle$

lemma *restriction-card-mon1*: $(v = u - - w) \longrightarrow |\text{restrict } u \ f \ c| \leq |\text{restrict } v \ f \ c|$
 $\langle \text{proof} \rangle$

lemma *restriction-card-mon2*: $(v = u - - w) \longrightarrow |\text{restrict } w \ f \ c| \leq |\text{restrict } v \ f \ c|$
 $\langle \text{proof} \rangle$

lemma *restriction-res-leq-two*: $|\text{restrict } v \ (\text{res } ts) \ c| \leq 2$
 $\langle \text{proof} \rangle$

lemma *restriction-clm-leq-one*: $|\text{restrict } v \ (\text{clm } ts) \ c| \leq 1$
 $\langle \text{proof} \rangle$

lemma *restriction-add-clm*:
 $(v = u - - w)$
 $\longrightarrow |\text{restrict } v \ (\text{clm } ts) \ c| = |\text{restrict } u \ (\text{clm } ts) \ c| + |\text{restrict } w \ (\text{clm } ts) \ c|$
 $\langle \text{proof} \rangle$

lemma *restriction-card-mon-trans*:
 $(v = v1 - - v2) \wedge (v2 = v3 - - v4) \wedge |\text{restrict } v3 \ f \ c| = 1 \longrightarrow |\text{restrict } v \ f \ c| \geq 1$
 $\langle \text{proof} \rangle$

lemma *restriction-card-somewhere-mon*:
 $(v = v1 \parallel v1) \wedge (v1 = v2 \parallel vr) \wedge (v2 = vu - - v3) \wedge (v3 = v' - - vd) \wedge |\text{restrict } v' \ f \ c| = 1$
 $\longrightarrow |\text{restrict } v \ f \ c| \geq 1$
 $\langle \text{proof} \rangle$

lemma *restrict-eq-lan-subs*:
 $|\text{restrict } v \ f \ c| = |\text{lan } v| \wedge (\text{restrict } v \ f \ c \sqsubseteq \text{lan } v) \longrightarrow \text{restrict } v \ f \ c = \text{lan } v$
 $\langle \text{proof} \rangle$

lemma *create-reservation-restrict-union*:
 $(ts - r(c) \longrightarrow ts')$
 $\longrightarrow \text{restrict } v \ (\text{res } ts') \ c = \text{restrict } v \ (\text{res } ts) \ c \sqcup \text{restrict } v \ (\text{clm } ts) \ c$
 $\langle \text{proof} \rangle$

lemma *switch-restrict-stable*: $(v = c > u) \longrightarrow \text{restrict } v \ f \ d = \text{restrict } u \ f \ d$
 $\langle \text{proof} \rangle$

end
end

7 Move a View according to Difference between Traffic Snapshots

In this section, we define a function to move a view according to the changes between two traffic snapshots. The intuition is that the view moves with the same speed as its owner. That is, if we move a view v from ts to ts' , we shift the extension of the view by the difference in the position of the owner of v .

```
theory Move
imports Traffic Views
begin
```

```
context traffic
begin
```

```
definition move::traffic  $\Rightarrow$  traffic  $\Rightarrow$  view  $\Rightarrow$  view
```

```
where
```

```
  move ts ts' v = ( $\parallel$  ext = shift (ext v) ((pos ts' (own v)) - pos ts (own v)),
                  lan = lan v,
                  own = own v  $\parallel$ )
```

```
lemma move-keeps-length:  $\parallel$  ext v  $\parallel$  =  $\parallel$  ext (move ts ts' v)  $\parallel$ 
   $\langle$ proof $\rangle$ 
```

```
lemma move-keeps-lanes: lan v = lan (move ts ts' v)  $\langle$ proof $\rangle$ 
```

```
lemma move-keeps-owner: own v = own (move ts ts' v)  $\langle$ proof $\rangle$ 
```

```
lemma move-nothing : move ts ts v = v  $\langle$ proof $\rangle$ 
```

```
lemma move-trans:
```

```
  (ts  $\Rightarrow$  ts')  $\wedge$  (ts'  $\Rightarrow$  ts'')  $\longrightarrow$  move ts' ts'' (move ts ts' v) = move ts ts'' v
   $\langle$ proof $\rangle$ 
```

```
lemma move-stability-res: (ts - r(c)  $\rightarrow$  ts')  $\longrightarrow$  move ts ts' v = v
```

```
and move-stability-clm: (ts - c(c,n)  $\rightarrow$  ts')  $\longrightarrow$  move ts ts' v = v
```

```
and move-stability-wdr: (ts - wdr(c,n)  $\rightarrow$  ts')  $\longrightarrow$  move ts ts' v = v
```

```
and move-stability-wdc: (ts - wdc(c)  $\rightarrow$  ts')  $\longrightarrow$  move ts ts' v = v
```

```
   $\langle$ proof $\rangle$ 
```

```
end
```

```
end
```

8 Sensors for Cars

This section presents the abstract definition of a function determining the sensor capabilities of cars. Such a function takes a car e , a traffic snapshot

ts and another car c , and returns the length of c as perceived by e at the situation determined by ts . The only restriction we impose is that this length is always greater than zero.

With such a function, we define a derived notion of the *space* the car c occupies as perceived by e . However, this does not define the lanes c occupies, but only a continuous interval. The lanes occupied by c are given by the reservation and claim functions of the traffic snapshot ts .

theory *Sensors*

imports *Traffic Views*

begin

locale *sensors* = *traffic* + *view* +

fixes *sensors*::(*cars*) \Rightarrow *traffic* \Rightarrow (*cars*) \Rightarrow *real*

assumes *sensors-ge*:(*sensors* e *ts* c) > 0

begin

definition *space* :: *traffic* \Rightarrow *view* \Rightarrow *cars* \Rightarrow *real-int*

where *space* *ts* v $c \equiv$ *Abs-real-int* (*pos* *ts* c , *pos* *ts* c + *sensors* (*own* v) *ts* c)

lemma *left-space*: *left* (*space* *ts* v c) = *pos* *ts* c

<proof>

lemma *right-space*: *right* (*space* *ts* v c) = *pos* *ts* c + *sensors* (*own* v) *ts* c

<proof>

lemma *space-nonempty*:*left* (*space* *ts* v c) $<$ *right* (*space* *ts* v c)

<proof>

end

end

9 Visible Length of Cars with Perfect Sensors

Given a sensor function, we can define the length of a car c as perceived by the owner of a view v . This length is restricted by the size of the extension of the view v , but always given by a continuous interval, which may possibly be degenerate (i.e., a point-interval).

The lemmas connect the end-points of the perceived length with the end-points of the current view. Furthermore, they show how the chopping and subview relations affect the perceived length of a car.

theory *Length*

imports *Sensors*

begin

context *sensors*

begin

definition *len*:: *view* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real-int*

where *len-def* : *len* *v* (*ts*) *c* ==
 if (*left* (*space* *ts* *v* *c*) > *right* (*ext* *v*))
 then *Abs-real-int* (*right* (*ext* *v*), *right* (*ext* *v*))
 else
 if (*right* (*space* *ts* *v* *c*) < *left* (*ext* *v*))
 then *Abs-real-int* (*left* (*ext* *v*), *left* (*ext* *v*))
 else
Abs-real-int (*max* (*left* (*ext* *v*)) (*left* (*space* *ts* *v* *c*)),
min (*right* (*ext* *v*)) (*right* (*space* *ts* *v* *c*)))

lemma *len-left*: *left* ((*len* *v* *ts*) *c*) \geq *left* (*ext* *v*)
 \langle *proof* \rangle

lemma *len-right*: *right* ((*len* *v* *ts*) *c*) \leq *right* (*ext* *v*)
 \langle *proof* \rangle

lemma *len-sub-int*: *len* *v* *ts* *c* \leq *ext* *v*
 \langle *proof* \rangle

lemma *len-space-left*:
left (*space* *ts* *v* *c*) \leq *right* (*ext* *v*) \longrightarrow *left* (*len* *v* *ts* *c*) \geq *left* (*space* *ts* *v* *c*)
 \langle *proof* \rangle

lemma *len-space-right*:
right (*space* *ts* *v* *c*) \geq *left* (*ext* *v*) \longrightarrow *right* (*len* *v* *ts* *c*) \leq *right* (*space* *ts* *v* *c*)
 \langle *proof* \rangle

lemma *len-hchop-left-right-border*:
(*len* *v* *ts* *c* = *ext* *v*) \wedge (*v*=*v1*||*v2*) \longrightarrow (*right* (*len* *v1* *ts* *c*) = *right* (*ext* *v1*))
 \langle *proof* \rangle

lemma *len-hchop-left-left-border*:
((*len* *v* *ts*) *c* = *ext* *v*) \wedge (*v*=*v1*||*v2*) \longrightarrow (*left* ((*len* *v1* *ts*) *c*) = *left* (*ext* *v1*))
 \langle *proof* \rangle

lemma *len-view-hchop-left*:
((*len* *v* *ts*) *c* = *ext* *v*) \wedge (*v*=*v1*||*v2*) \longrightarrow ((*len* *v1* *ts*) *c* = *ext* *v1*)
 \langle *proof* \rangle

lemma *len-hchop-right-left-border*:
((*len* *v* *ts*) *c* = *ext* *v*) \wedge (*v*=*v1*||*v2*) \longrightarrow (*left* ((*len* *v2* *ts*) *c*) = *left* (*ext* *v2*))
 \langle *proof* \rangle

lemma *len-hchop-right-right-border*:
((*len* *v* *ts*) *c* = *ext* *v*) \wedge (*v*=*v1*||*v2*) \longrightarrow (*right* ((*len* *v2* *ts*) *c*) = *right* (*ext* *v2*))

$\langle proof \rangle$

lemma *len-view-hchop-right*:

$$((len\ v\ ts)\ c = ext\ v) \wedge (v=v1 \parallel v2) \longrightarrow ((len\ v2\ ts)\ c = ext\ v2)$$

$\langle proof \rangle$

lemma *len-compose-hchop*:

$$(v=v1 \parallel v2) \wedge (len\ v1\ (ts)\ c = ext\ v1) \wedge (len\ v2\ (ts)\ c = ext\ v2) \\ \longrightarrow (len\ v\ (ts)\ c = ext\ v)$$

$\langle proof \rangle$

lemma *len-stable*: $(v=v1 \dashv\dashv v2) \longrightarrow len\ v1\ ts\ c = len\ v2\ ts\ c$

$\langle proof \rangle$

lemma *len-empty-on-subview1*:

$$\|len\ v\ (ts)\ c\| = 0 \wedge (v=v1 \parallel v2) \longrightarrow \|len\ v1\ (ts)\ c\| = 0$$

$\langle proof \rangle$

lemma *len-empty-on-subview2*:

$$\|len\ v\ ts\ c\| = 0 \wedge (v=v1 \parallel v2) \longrightarrow \|len\ v2\ ts\ c\| = 0$$

$\langle proof \rangle$

lemma *len-hchop-add*:

$$(v=v1 \parallel v2) \longrightarrow \|len\ v\ ts\ c\| = \|len\ v1\ ts\ c\| + \|len\ v2\ ts\ c\|$$

$\langle proof \rangle$

lemma *len-non-empty-inside*:

$$\|len\ v\ (ts)\ c\| > 0$$

$$\longrightarrow left\ (space\ ts\ v\ c) < right\ (ext\ v) \wedge right\ (space\ ts\ v\ c) > left\ (ext\ v)$$

$\langle proof \rangle$

lemma *len-fills-subview*:

$$\|len\ v\ ts\ c\| > 0$$

$$\longrightarrow (\exists\ v1\ v2\ v3\ v'. (v=v1 \parallel v2) \wedge (v2=v' \parallel v3) \wedge len\ v'\ ts\ c = ext\ v' \wedge \\ \|len\ v'\ ts\ c\| = \|len\ v\ ts\ c\|)$$

$\langle proof \rangle$

lemma *ext-eq-len-eq*:

$$ext\ v = ext\ v' \wedge own\ v = own\ v' \longrightarrow len\ v\ ts\ c = len\ v'\ ts\ c$$

$\langle proof \rangle$

lemma *len-stable-down*: $(v=v1 \dashv\dashv v2) \longrightarrow len\ v\ ts\ c = len\ v1\ ts\ c$

$\langle proof \rangle$

lemma *len-stable-up*: $(v=v1 \dashv\dashv v2) \longrightarrow len\ v\ ts\ c = len\ v2\ ts\ c$

$\langle proof \rangle$

lemma *len-empty-subview*: $\|len\ v\ ts\ c\| = 0 \wedge (v' \leq v) \longrightarrow \|len\ v'\ ts\ c\| = 0$

$\langle proof \rangle$

lemma *view-leq-len-leq*: $(ext\ v \leq ext\ v') \wedge (own\ v = own\ v') \wedge \|len\ v\ ts\ c\| > 0$
 $\longrightarrow len\ v\ ts\ c \leq len\ v'\ ts\ c$

$\langle proof \rangle$

end
end

10 Basic HMLSL

In this section, we define the basic formulas of HMLSL. All of these basic formulas and theorems are independent of the choice of sensor function. However, they show how the general operators (chop, changes in perspective, atomic formulas) work.

theory *HMLSL*
imports *Restriction Move Length*
begin

10.1 Syntax of Basic HMLSL

Formulas are functions associating a traffic snapshot and a view with a Boolean value.

type-synonym $\sigma = traffic \Rightarrow view \Rightarrow bool$

locale *hmlsl* = *restriction+*
fixes *sensors::cars* $\Rightarrow traffic \Rightarrow cars \Rightarrow real$
assumes *sensors-ge*: $(sensors\ e\ ts\ c) > 0$ **begin**
end

sublocale *hmlsl* < *sensors*
 $\langle proof \rangle$

context *hmlsl*
begin

All formulas are defined as abbreviations. As a consequence, proofs will directly refer to the semantics of HMLSL, i.e., traffic snapshots and views.

The first-order operators are direct translations into HOL operators.

abbreviation *mtrue* :: $\sigma \rightarrow \langle \top \rangle$
where $\top \equiv \lambda ts\ w. True$
abbreviation *mfalse* :: $\sigma \rightarrow \langle \perp \rangle$
where $\perp \equiv \lambda ts\ w. False$
abbreviation *mnot* :: $\sigma \Rightarrow \sigma \rightarrow \langle \neg \rightarrow [52]53 \rangle$
where $\neg\varphi \equiv \lambda ts\ w. \neg\varphi(ts)(w)$
abbreviation *mnegpred* :: $(cars \Rightarrow \sigma) \Rightarrow (cars \Rightarrow \sigma) \rightarrow \langle \neg \rightarrow [52]53 \rangle$

where $\neg\Phi \equiv \lambda x. \lambda ts w. \neg\Phi(x)(ts)(w)$
abbreviation *mand* $:: \sigma \Rightarrow \sigma \Rightarrow \sigma$ (**infixr** $\langle \wedge \rangle$ 51)
where $\varphi \wedge \psi \equiv \lambda ts w. \varphi(ts)(w) \wedge \psi(ts)(w)$
abbreviation *mor* $:: \sigma \Rightarrow \sigma \Rightarrow \sigma$ (**infix** $\langle \vee \rangle$ 50)
where $\varphi \vee \psi \equiv \lambda ts w. \varphi(ts)(w) \vee \psi(ts)(w)$
abbreviation *mimp* $:: \sigma \Rightarrow \sigma \Rightarrow \sigma$ (**infixr** $\langle \rightarrow \rangle$ 49)
where $\varphi \rightarrow \psi \equiv \lambda ts w. \varphi(ts)(w) \rightarrow \psi(ts)(w)$
abbreviation *mequ* $:: \sigma \Rightarrow \sigma \Rightarrow \sigma$ (**infixr** $\langle \leftrightarrow \rangle$ 48)
where $\varphi \leftrightarrow \psi \equiv \lambda ts w. \varphi(ts)(w) \leftrightarrow \psi(ts)(w)$
abbreviation *mforall* $:: ('a \Rightarrow \sigma) \Rightarrow \sigma$ ($\langle \forall \rangle$)
where $\forall \Phi \equiv \lambda ts w. \forall x. \Phi(x)(ts)(w)$
abbreviation *mforallB* $:: ('a \Rightarrow \sigma) \Rightarrow \sigma$ (**binder** $\langle \forall \rangle$ [8] 9)
where $\forall x. \varphi(x) \equiv \forall \varphi$
abbreviation *mexists* $:: ('a \Rightarrow \sigma) \Rightarrow \sigma$ ($\langle \exists \rangle$)
where $\exists \Phi \equiv \lambda ts w. \exists x. \Phi(x)(ts)(w)$
abbreviation *mexistsB* $:: (('a) \Rightarrow \sigma) \Rightarrow \sigma$ (**binder** $\langle \exists \rangle$ [8] 9)
where $\exists x. \varphi(x) \equiv \exists \varphi$
abbreviation *meq* $:: 'a \Rightarrow 'a \Rightarrow \sigma$ (**infixr** $\langle = \rangle$ 60) — Equality
where $x = y \equiv \lambda ts w. x = y$
abbreviation *mgeq* $:: ('a :: \text{ord}) \Rightarrow 'a \Rightarrow \sigma$ (**infix** $\langle \geq \rangle$ 60)
where $x \geq y \equiv \lambda ts w. x \geq y$
abbreviation *mge* $:: ('a :: \text{ord}) \Rightarrow 'a \Rightarrow \sigma$ (**infix** $\langle > \rangle$ 60)
where $x > y \equiv \lambda ts w. x > y$

For the spatial modalities, we use the chopping operations defined on views. Observe that our chop modalities are existential.

abbreviation *hchop* $:: \sigma \Rightarrow \sigma \Rightarrow \sigma$ (**infixr** $\langle \frown \rangle$ 53)
where $\varphi \frown \psi \equiv \lambda ts w. \exists v u. (w = v \parallel u) \wedge \varphi(ts)(v) \wedge \psi(ts)(u)$
abbreviation *vchop* $:: \sigma \Rightarrow \sigma \Rightarrow \sigma$ (**infixr** $\langle \smile \rangle$ 53)
where $\varphi \smile \psi \equiv \lambda ts w. \exists v u. (w = v \dashv\!\! \dashv u) \wedge \varphi(ts)(v) \wedge \psi(ts)(u)$
abbreviation *somewhere* $:: \sigma \Rightarrow \sigma$ ($\langle \langle - \rangle \rangle$ 55)
where $\langle \varphi \rangle \equiv \top \frown (\top \smile \varphi \smile \top) \frown \top$
abbreviation *everywhere* $:: \sigma \Rightarrow \sigma$ ($\langle [-] \rangle$ 55)
where $[\varphi] \equiv \neg \langle \neg \varphi \rangle$

To change the perspective of a view, we use an operator in the fashion of Hybrid Logic.

abbreviation *at* $:: \text{cars} \Rightarrow \sigma \Rightarrow \sigma$ ($\langle @ - \rightarrow \rangle$ 56)
where $@_c \varphi \equiv \lambda ts w. \forall v'. (w = c > v') \rightarrow \varphi(ts)(v')$

The behavioural modalities are defined as usual modal box-like modalities, where the accessibility relations are given by the different types of transitions between traffic snapshots.

abbreviation *res-box* $:: \text{cars} \Rightarrow \sigma \Rightarrow \sigma$ ($\langle \Box r'(-) \rightarrow \rangle$ 55)
where $\Box r(c) \varphi \equiv \lambda ts w. \forall ts'. (ts \text{--} r(c) \rightarrow ts') \rightarrow \varphi(ts')(w)$
abbreviation *clm-box* $:: \text{cars} \Rightarrow \sigma \Rightarrow \sigma$ ($\langle \Box c'(-) \rightarrow \rangle$ 55)
where $\Box c(c) \varphi \equiv \lambda ts w. \forall ts' n. (ts \text{--} c(c, n) \rightarrow ts') \rightarrow \varphi(ts')(w)$
abbreviation *wdres-box* $:: \text{cars} \Rightarrow \sigma \Rightarrow \sigma$ ($\langle \Box wdr'(-) \rightarrow \rangle$ 55)

where $\Box wdr(c) \varphi \equiv \lambda ts w. \forall ts' n. (ts - wdr(c, n) \rightarrow ts') \rightarrow \varphi(ts')(w)$
abbreviation $wdclm\text{-}box::cars \Rightarrow \sigma \Rightarrow \sigma (\langle \Box wdc'(-) \rangle \rightarrow 55)$
where $\Box wdc(c) \varphi \equiv \lambda ts w. \forall ts'. (ts - wdc(c) \rightarrow ts') \rightarrow \varphi(ts')(w)$
abbreviation $time\text{-}box::\sigma \Rightarrow \sigma (\langle \Box \tau \rangle \rightarrow 55)$
where $\Box \tau \varphi \equiv \lambda ts w. \forall ts'. (ts \rightsquigarrow ts') \rightarrow \varphi(ts')(move\ ts\ ts'\ w)$
abbreviation $globally::\sigma \Rightarrow \sigma (\langle \mathbf{G} \rangle \rightarrow 55)$
where $\mathbf{G} \varphi \equiv \lambda ts w. \forall ts'. (ts \Rightarrow ts') \rightarrow \varphi(ts')(move\ ts\ ts'\ w)$

The spatial atoms to refer to reservations, claims and free space are direct translations of the original definitions of MLSL [2] into the Isabelle implementation.

abbreviation $re:: cars \Rightarrow \sigma (\langle re'(-) \rangle \rightarrow 70)$

where

$$\begin{aligned}
 re(c) &\equiv \lambda ts v. \|ext\ v\| > 0 \wedge len\ v\ ts\ c = ext\ v \wedge \\
 &\quad restrict\ v\ (res\ ts)\ c = lan\ v \wedge |lan\ v| = 1
 \end{aligned}$$

abbreviation $cl:: cars \Rightarrow \sigma (\langle cl'(-) \rangle \rightarrow 70)$

where

$$\begin{aligned}
 cl(c) &\equiv \lambda ts v. \|ext\ v\| > 0 \wedge len\ v\ ts\ c = ext\ v \wedge \\
 &\quad restrict\ v\ (clm\ ts)\ c = lan\ v \wedge |lan\ v| = 1
 \end{aligned}$$

abbreviation $free:: \sigma (\langle free \rangle)$

where

$$\begin{aligned}
 free &\equiv \lambda ts v. \|ext\ v\| > 0 \wedge |lan\ v| = 1 \wedge \\
 &\quad (\forall c. \|len\ v\ ts\ c\| = 0 \vee \\
 &\quad (restrict\ v\ (clm\ ts)\ c = \emptyset \wedge restrict\ v\ (res\ ts)\ c = \emptyset))
 \end{aligned}$$

Even though we do not need them for the subsequent proofs of safety, we define ways to measure the number of lanes (width) and the size of the extension (length) of a view. This allows us to connect the atomic formulas for reservations and claims with the atom denoting free space [5].

abbreviation $width\text{-}eq::nat \Rightarrow \sigma (\langle \omega = - \rangle \rightarrow 60)$

where $\omega = n \equiv \lambda ts v. |lan\ v| = n$

abbreviation $width\text{-}geq::nat \Rightarrow \sigma (\langle \omega \geq - \rangle \rightarrow 60)$

where $\omega \geq n \equiv \lambda ts v. |lan\ v| \geq n$

abbreviation $width\text{-}ge::nat \Rightarrow \sigma (\langle \omega > - \rangle \rightarrow 60)$

where $\omega > n \equiv (\omega = n+1) \smile \top$

abbreviation $length\text{-}eq::real \Rightarrow \sigma (\langle l = - \rangle \rightarrow 60)$

where $l = r \equiv \lambda ts v. \|ext\ v\| = r$

abbreviation $length\text{-}ge::real \Rightarrow \sigma (\langle l > - \rangle \rightarrow 60)$

where $l > r \equiv \lambda ts v. \|ext\ v\| > r$

abbreviation $length\text{-}geq::real \Rightarrow \sigma (\langle l \geq - \rangle \rightarrow 60)$

where $l \geq r \equiv (l = r) \vee (l > r)$

For convenience, we use abbreviations for the validity and satisfiability of formulas. While the former gives a nice way to express theorems, the latter is useful within proofs.

abbreviation *valid* :: $\sigma \Rightarrow \text{bool}$ ($\langle \models \rightarrow 10 \rangle$)
where $\models \varphi \equiv \forall ts. \forall v. \varphi(ts)(v)$

abbreviation *satisfies* :: $\text{traffic} \Rightarrow \text{view} \Rightarrow \sigma \Rightarrow \text{bool}$ ($\langle -, - \models \rightarrow 10 \rangle$)
where $ts, v \models \varphi \equiv \varphi(ts)(v)$

10.2 Theorems about Basic HMLSL

lemma *hchop-weaken1*: $\models \varphi \rightarrow (\varphi \frown \top)$
 $\langle \text{proof} \rangle$

lemma *hchop-weaken2*: $\models \varphi \rightarrow (\top \frown \varphi)$
 $\langle \text{proof} \rangle$

lemma *hchop-weaken*: $\models \varphi \rightarrow (\top \frown \varphi \frown \top)$
 $\langle \text{proof} \rangle$

lemma *hchop-neg1*: $\models \neg(\varphi \frown \top) \rightarrow ((\neg \varphi) \frown \top)$
 $\langle \text{proof} \rangle$

lemma *hchop-neg2*: $\models \neg(\top \frown \varphi) \rightarrow (\top \frown \neg \varphi)$
 $\langle \text{proof} \rangle$

lemma *hchop-disj-distr1*: $\models ((\varphi \frown (\psi \vee \chi)) \leftrightarrow ((\varphi \frown \psi) \vee (\varphi \frown \chi)))$
 $\langle \text{proof} \rangle$

lemma *hchop-disj-distr2*: $\models (((\psi \vee \chi) \frown \varphi) \leftrightarrow ((\psi \frown \varphi) \vee (\chi \frown \varphi)))$
 $\langle \text{proof} \rangle$

lemma *hchop-assoc*: $\models \varphi \frown (\psi \frown \chi) \leftrightarrow (\varphi \frown \psi) \frown \chi$
 $\langle \text{proof} \rangle$

lemma *v-chop-weaken1*: $\models (\varphi \rightarrow (\varphi \smile \top))$
 $\langle \text{proof} \rangle$

lemma *v-chop-weaken2*: $\models (\varphi \rightarrow (\top \smile \varphi))$
 $\langle \text{proof} \rangle$

lemma *v-chop-assoc*: $\models (\varphi \smile (\psi \smile \chi)) \leftrightarrow ((\varphi \smile \psi) \smile \chi)$
 $\langle \text{proof} \rangle$

lemma *vchop-disj-distr1*: $\models ((\varphi \smile (\psi \vee \chi)) \leftrightarrow ((\varphi \smile \psi) \vee (\varphi \smile \chi)))$
 $\langle \text{proof} \rangle$

lemma *vchop-disj-distr2*: $\models (((\psi \vee \chi) \smile \varphi) \leftrightarrow ((\psi \smile \varphi) \vee (\chi \smile \varphi)))$
 $\langle \text{proof} \rangle$

lemma *at-exists* : $\models \varphi \rightarrow (\exists c. @c \varphi)$

<proof>

lemma *at-conj-dist* : $\models (@c (\varphi \wedge \psi)) \leftrightarrow ((@c \varphi) \wedge (@c \psi))$

<proof>

lemma *at-disj-dist* : $\models (@c (\varphi \vee \psi)) \leftrightarrow ((@c \varphi) \vee (@c \psi))$

<proof>

lemma *at-hchop-dist1* : $\models (@c (\varphi \frown \psi)) \rightarrow ((@c \varphi) \frown (@c \psi))$

<proof>

lemma *at-hchop-dist2* : $\models ((@c \varphi) \frown (@c \psi)) \rightarrow (@c (\varphi \frown \psi))$

<proof>

lemma *at-hchop-dist* : $\models ((@c \varphi) \frown (@c \psi)) \leftrightarrow (@c (\varphi \frown \psi))$

<proof>

lemma *at-vchop-dist1* : $\models (@c (\varphi \smile \psi)) \rightarrow ((@c \varphi) \smile (@c \psi))$

<proof>

lemma *at-vchop-dist2* : $\models ((@c \varphi) \smile (@c \psi)) \rightarrow (@c (\varphi \smile \psi))$

<proof>

lemma *at-vchop-dist* : $\models ((@c \varphi) \smile (@c \psi)) \leftrightarrow (@c (\varphi \smile \psi))$

<proof>

lemma *at-eq* : $\models (@e c = d) \leftrightarrow (c = d)$

<proof>

lemma *at-neg1* : $\models (@c \neg \varphi) \rightarrow \neg (@c \varphi)$

<proof>

lemma *at-neg2* : $\models \neg (@c \varphi) \rightarrow (@c \neg \varphi)$

<proof>

lemma *at-neg* : $\models (@c (\neg \varphi)) \leftrightarrow \neg (@c \varphi)$

<proof>

lemma *at-neg'ts,v* : $\models \neg (@c \varphi) \leftrightarrow (@c (\neg \varphi))$ *<proof>*

lemma *at-neg-neg1* : $\models (@c \varphi) \rightarrow \neg (@c \neg \varphi)$

<proof>

lemma *at-neg-neg2* : $\models \neg (@c \neg \varphi) \rightarrow (@c \varphi)$

<proof>

lemma *at-neg-neg* : $\models (@c \varphi) \leftrightarrow \neg (@c \neg \varphi)$

$\langle proof \rangle$

lemma *globally-all-iff*: $\models (\mathbf{G}(\forall c. \varphi)) \leftrightarrow (\forall c. (\mathbf{G} \varphi))$ $\langle proof \rangle$

lemma *globally-all-iff':ts,v*: $\models (\mathbf{G}(\forall c. \varphi)) \leftrightarrow (\forall c. (\mathbf{G} \varphi))$ $\langle proof \rangle$

lemma *globally-reft*: $\models (\mathbf{G} \varphi) \rightarrow \varphi$
 $\langle proof \rangle$

lemma *globally-4*: $\models (\mathbf{G} \varphi) \rightarrow \mathbf{G} \mathbf{G} \varphi$
 $\langle proof \rangle$

lemma *spatial-weaken*: $\models (\varphi \rightarrow \langle \varphi \rangle)$
 $\langle proof \rangle$

lemma *spatial-weaken2*: $\models (\varphi \rightarrow \psi) \rightarrow (\varphi \rightarrow \langle \psi \rangle)$
 $\langle proof \rangle$

lemma *somewhere-distr*: $\models \langle \varphi \vee \psi \rangle \leftrightarrow \langle \varphi \rangle \vee \langle \psi \rangle$
 $\langle proof \rangle$

lemma *somewhere-and*: $\models \langle \varphi \wedge \psi \rangle \rightarrow \langle \varphi \rangle \wedge \langle \psi \rangle$
 $\langle proof \rangle$

lemma *somewhere-and-or-distr*: $\models (\langle \chi \wedge (\varphi \vee \psi) \rangle \leftrightarrow \langle \chi \wedge \varphi \rangle \vee \langle \chi \wedge \psi \rangle)$
 $\langle proof \rangle$

lemma *width-add1*: $\models ((\omega = x) \smile (\omega = y) \rightarrow \omega = x+y)$
 $\langle proof \rangle$

lemma *width-add2*: $\models ((\omega = x+y) \rightarrow (\omega = x) \smile \omega = y)$
 $\langle proof \rangle$

lemma *width-hchop-stable*: $\models ((\omega = x) \leftrightarrow ((\omega = x) \smile (\omega = x)))$
 $\langle proof \rangle$

lemma *length-geq-zero*: $\models (1 \geq 0)$
 $\langle proof \rangle$

lemma *length-split*: $\models ((1 > 0) \rightarrow (1 > 0) \smile (1 > 0))$
 $\langle proof \rangle$

lemma *length-meld*: $\models ((1 > 0) \smile (1 > 0) \rightarrow (1 > 0))$
 $\langle proof \rangle$

lemma *length-dense*: $\models ((1 > 0) \leftrightarrow (1 > 0) \smile (1 > 0))$
 $\langle proof \rangle$

lemma *length-add1*: $\models ((1 = x) \smile (1 = y)) \rightarrow (1 = x+y)$

$\langle proof \rangle$

lemma *length-add2*: $\models (x \geq 0 \wedge y \geq 0) \rightarrow ((l=x+y) \rightarrow ((l=x) \frown (l=y)))$
 $\langle proof \rangle$

lemma *length-add*: $\models (x \geq 0 \wedge y \geq 0) \rightarrow ((l=x+y) \leftrightarrow ((l=x) \frown (l=y)))$
 $\langle proof \rangle$

lemma *length-vclop-stable*: $\models (l = x) \leftrightarrow ((l = x) \smile (l = x))$
 $\langle proof \rangle$

lemma *res-ge-zero*: $\models (re(c) \rightarrow l > 0)$
 $\langle proof \rangle$

lemma *clm-ge-zero*: $\models (cl(c) \rightarrow l > 0)$
 $\langle proof \rangle$

lemma *free-ge-zero*: $\models free \rightarrow l > 0$
 $\langle proof \rangle$

lemma *width-res*: $\models (re(c) \rightarrow \omega = 1)$
 $\langle proof \rangle$

lemma *width-clm*: $\models (cl(c) \rightarrow \omega = 1)$
 $\langle proof \rangle$

lemma *width-free*: $\models (free \rightarrow \omega = 1)$
 $\langle proof \rangle$

lemma *width-somewhere-res*: $\models \langle re(c) \rangle \rightarrow (\omega \geq 1)$
 $\langle proof \rangle$

lemma *clm-disj-res*: $\models \neg \langle cl(c) \wedge re(c) \rangle$
 $\langle proof \rangle$

lemma *width-ge*: $\models (\omega > 0) \rightarrow (\exists x. (\omega = x) \wedge (x > 0))$
 $\langle proof \rangle$

lemma *two-res-width*: $\models ((re(c) \smile re(c)) \rightarrow \omega = 2)$
 $\langle proof \rangle$

lemma *res-at-most-two*: $\models \neg (re(c) \smile re(c) \smile re(c))$
 $\langle proof \rangle$

lemma *res-at-most-two2*: $\models \neg \langle re(c) \smile re(c) \smile re(c) \rangle$
 $\langle proof \rangle$

lemma *res-at-most-somewhere*: $\models \neg \langle re(c) \rangle \smile \langle re(c) \rangle \smile \langle re(c) \rangle$

$\langle proof \rangle$

lemma *res-adj*: $\models \neg (re(c) \smile (\omega > 0) \smile re(c))$
 $\langle proof \rangle$

lemma *clm-sing*: $\models \neg (cl(c) \smile cl(c))$
 $\langle proof \rangle$

lemma *clm-sing-somewhere*: $\models \neg \langle cl(c) \smile cl(c) \rangle$
 $\langle proof \rangle$

lemma *clm-sing-not-interrupted*: $\models \neg (cl(c) \smile \top \smile cl(c))$
 $\langle proof \rangle$

lemma *clm-sing-somewhere2*: $\models \neg (\top \smile cl(c) \smile \top \smile cl(c) \smile \top)$
 $\langle proof \rangle$

lemma *clm-sing-somewhere3*: $\models \neg \langle (\top \smile cl(c) \smile \top \smile cl(c) \smile \top) \rangle$
 $\langle proof \rangle$

lemma *clm-at-most-somewhere*: $\models \neg (\langle cl(c) \rangle \smile \langle cl(c) \rangle)$
 $\langle proof \rangle$

lemma *res-decompose*: $\models (re(c) \rightarrow re(c) \frown re(c))$
 $\langle proof \rangle$

lemma *res-compose*: $\models (re(c) \frown re(c) \rightarrow re(c))$
 $\langle proof \rangle$

lemma *res-dense*: $\models re(c) \leftrightarrow re(c) \frown re(c)$
 $\langle proof \rangle$

lemma *res-continuous*: $\models (re(c) \rightarrow (\neg (\top \frown (\neg re(c) \wedge 1 > 0) \frown \top)))$
 $\langle proof \rangle$

lemma *no-clm-before-res*: $\models \neg (cl(c) \frown re(c))$
 $\langle proof \rangle$

lemma *no-clm-before-res2*: $\models \neg (cl(c) \frown \top \frown re(c))$
 $\langle proof \rangle$

lemma *clm-decompose*: $\models (cl(c) \rightarrow cl(c) \frown cl(c))$
 $\langle proof \rangle$

lemma *clm-compose*: $\models (cl(c) \frown cl(c) \rightarrow cl(c))$

$\langle proof \rangle$

lemma *clm-dense*: $\models cl(c) \leftrightarrow cl(c) \frown cl(c)$
 $\langle proof \rangle$

lemma *clm-continuous*: $\models (cl(c) \rightarrow (\neg (\top \frown (\neg cl(c) \wedge \mathbf{l} > 0) \frown \top)))$
 $\langle proof \rangle$

lemma *res-not-free*: $\models (\exists c. re(c) \rightarrow \neg free)$
 $\langle proof \rangle$

lemma *clm-not-free*: $\models (\exists c. cl(c) \rightarrow \neg free)$
 $\langle proof \rangle$

lemma *free-no-res*: $\models (free \rightarrow \neg (\exists c. re(c)))$
 $\langle proof \rangle$

lemma *free-no-clm*: $\models (free \rightarrow \neg (\exists c. cl(c)))$
 $\langle proof \rangle$

lemma *free-decompose*: $\models free \rightarrow (free \frown free)$
 $\langle proof \rangle$

lemma *free-compose*: $\models (free \frown free) \rightarrow free$
 $\langle proof \rangle$

lemma *free-dense*: $\models free \leftrightarrow (free \frown free)$
 $\langle proof \rangle$

lemma *free-dense2*: $\models free \rightarrow \top \frown free \frown \top$
 $\langle proof \rangle$

The next lemmas show the connection between the spatial. In particular, if the view consists of one lane and a non-zero extension, where neither a reservation nor a car resides, the view satisfies free (and vice versa).

lemma *no-cars-means-free*:
 $\models ((\mathbf{l} > 0) \wedge (\omega = 1) \wedge (\forall c. \neg (\top \frown (cl(c) \vee re(c)) \frown \top))) \rightarrow free$
 $\langle proof \rangle$

lemma *free-means-no-cars*:
 $\models free \rightarrow ((\mathbf{l} > 0) \wedge (\omega = 1) \wedge (\forall c. \neg (\top \frown (cl(c) \vee re(c)) \frown \top)))$
 $\langle proof \rangle$

lemma *free-eq-no-cars*:
 $\models free \leftrightarrow ((\mathbf{l} > 0) \wedge (\omega = 1) \wedge (\forall c. \neg (\top \frown (cl(c) \vee re(c)) \frown \top)))$
 $\langle proof \rangle$

lemma *free-nowhere-res*: $\models \text{free} \rightarrow \neg(\top \frown (\text{re}(c)) \frown \top)$
 $\langle \text{proof} \rangle$

lemma *two-res-not-res*: $\models ((\text{re}(c) \smile \text{re}(c)) \rightarrow \neg \text{re}(c))$
 $\langle \text{proof} \rangle$

lemma *two-clm-width*: $\models ((\text{cl}(c) \smile \text{cl}(c)) \rightarrow \omega = 2)$
 $\langle \text{proof} \rangle$

lemma *two-res-no-car*: $\models (\text{re}(c) \smile \text{re}(c)) \rightarrow \neg(\exists c. (\text{cl}(c) \vee \text{re}(c)))$
 $\langle \text{proof} \rangle$

lemma *two-lanes-no-car*: $\models (\neg \omega = 1) \rightarrow \neg(\exists c. (\text{cl}(c) \vee \text{re}(c)))$
 $\langle \text{proof} \rangle$

lemma *empty-no-car*: $\models (1 = 0) \rightarrow \neg(\exists c. (\text{cl}(c) \vee \text{re}(c)))$
 $\langle \text{proof} \rangle$

lemma *car-one-lane-non-empty*: $\models (\exists c. (\text{cl}(c) \vee \text{re}(c))) \rightarrow ((\omega = 1) \wedge (1 > 0))$
 $\langle \text{proof} \rangle$

lemma *one-lane-notfree*:
 $\models (\omega = 1) \wedge (1 > 0) \wedge (\neg \text{free}) \rightarrow ((\top \frown (\exists c. (\text{re}(c) \vee \text{cl}(c))) \frown \top)$
 $\langle \text{proof} \rangle$

lemma *one-lane-empty-or-car*:
 $\models (\omega = 1) \wedge (1 > 0) \rightarrow (\text{free} \vee (\top \frown (\exists c. (\text{re}(c) \vee \text{cl}(c))) \frown \top))$
 $\langle \text{proof} \rangle$

end
end

11 Perfect Sensors

This section contains an instantiations of the sensor function for "perfect sensors". That is, each car can perceive both the physical size as well as the braking distance of each other car.

theory *Perfect-Sensors*
imports *../Length*
begin

definition *perfect*::*cars* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real*
where *perfect* *e ts c* \equiv *traffic.physical-size ts c* + *traffic.braking-distance ts c*

locale *perfect-sensors* = *traffic+view*
begin

interpretation *perfect-sensors* : *sensors perfect* :: *cars* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real*
 $\langle \text{proof} \rangle$

notation *perfect-sensors.space* ($\langle space \rangle$)

notation *perfect-sensors.len* ($\langle len \rangle$)

With this sensor definition, we can show that the perceived length of a car is independent of the spatial transitions between traffic snapshots. The length may only change during evolutions, in particular if the car changes its dynamical behaviour.

lemma *create-reservation-length-stable*:

$(ts - r(d) \rightarrow ts') \longrightarrow len\ v\ ts\ c = len\ v\ ts'\ c$
 $\langle proof \rangle$

lemma *create-claim-length-stable*:

$(ts - c(d, n) \rightarrow ts') \longrightarrow len\ v\ ts\ c = len\ v\ ts'\ c$
 $\langle proof \rangle$

lemma *withdraw-reservation-length-stable*:

$(ts - wdr(d, n) \rightarrow ts') \longrightarrow len\ v\ ts\ c = len\ v\ ts'\ c$
 $\langle proof \rangle$

lemma *withdraw-claim-length-stable*:

$(ts - wdc(d) \rightarrow ts') \longrightarrow len\ v\ ts\ c = len\ v\ ts'\ c$
 $\langle proof \rangle$

The following lemma shows that the perceived length is independent from the owner of the view. That is, as long as two views consist of the same extension, the perceived length of each car is the same in both views.

lemma *all-own-ext-eq-len-eq*:

$ext\ v = ext\ v' \longrightarrow len\ v\ ts\ c = len\ v'\ ts\ c$
 $\langle proof \rangle$

Finally, switching the perspective of a view does not change the perceived length.

lemma *switch-length-stable*: $(v = d \triangleright v') \longrightarrow len\ v\ ts\ c = len\ v'\ ts\ c$
 $\langle proof \rangle$

end

end

12 HMLSL for Perfect Sensors

Within this section, we instantiate HMLSL for cars with perfect sensors.

theory *HMLSL-Perfect*

imports *../HMLSL Perfect-Sensors*

begin

locale *hmlsl-perfect* = *perfect-sensors* + *restriction*
begin

interpretation *hmlsl* : *hmlsl perfect* :: *cars* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real*
 $\langle proof \rangle$

notation *hmlsl.re* ($\langle re'(-) \rangle$)

notation *hmlsl.cl* ($\langle cl'(-) \rangle$)

notation *hmlsl.len* ($\langle len \rangle$)

The spatial atoms are independent of the perspective of the view. Hence we can prove several lemmas on the relation between the hybrid modality and the spatial atoms.

lemma *at-res1*: $\models (re(c)) \rightarrow (\forall d. @d re(c))$
 $\langle proof \rangle$

lemma *at-res2*: $\models (\forall d. @d re(c)) \rightarrow re(c)$
 $\langle proof \rangle$

lemma *at-res*: $\models re(c) \leftrightarrow (\forall d. @d re(c))$
 $\langle proof \rangle$

lemma *at-res-inst*: $\models (@d re(c)) \rightarrow re(c)$
 $\langle proof \rangle$

lemma *at-clm1*: $\models cl(c) \rightarrow (\forall d. @d cl(c))$
 $\langle proof \rangle$

lemma *at-clm2*: $\models (\forall d. @d cl(c)) \rightarrow cl(c)$
 $\langle proof \rangle$

lemma *at-clm*: $\models cl(c) \leftrightarrow (\forall d. @d cl(c))$
 $\langle proof \rangle$

lemma *at-clm-inst*: $\models (@d cl(c)) \rightarrow cl(c)$
 $\langle proof \rangle$

With the definition of sensors, we can also express how the spatial situation changes after the different transitions. In particular, we can prove lemmas corresponding to the activity and stability rules of the proof system for MLSL [5].

Observe that we were not able to prove these rules for basic HMLSL, since its generic sensor function allows for instantiations where the perceived length changes during spatial transitions.

lemma *backwards-res-act*:

$(ts \rightarrow r(c) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c) \vee cl(c))$
 $\langle proof \rangle$

lemma *backwards-res-act-somewhere*:

$$(ts - r(c) \rightarrow ts') \wedge (ts', v \models \langle re(c) \rangle) \longrightarrow (ts, v \models \langle re(c) \vee cl(c) \rangle)$$

<proof>

lemma *backwards-res-stab*:

$$(ts - r(d) \rightarrow ts') \wedge (d \neq c) \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$$

<proof>

lemma *backwards-c-res-stab*:

$$(ts - c(d, n) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$$

<proof>

lemma *backwards-wdc-res-stab*:

$$(ts - wdc(d) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$$

<proof>

lemma *backwards-wdr-res-stab*:

$$(ts - wdr(d, n) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$$

<proof>

We now proceed to prove the *reservation lemma*, which was crucial in the manual safety proof [2].

lemma *reservation1*: $\models (re(c) \vee cl(c)) \rightarrow \Box r(c) re(c)$

<proof>

lemma *reservation2*: $\models (\Box r(c) re(c)) \rightarrow (re(c) \vee cl(c))$

<proof>

lemma *reservation*: $\models (\Box r(c) re(c)) \leftrightarrow (re(c) \vee cl(c))$

<proof>

end

end

13 Safety for Cars with Perfect Sensors

This section contains the definition of requirements for lane change and distance controllers for cars, with the assumption of perfect sensors. Using these definitions, we show that safety is an invariant along all possible behaviour of cars.

theory *Safety-Perfect*

imports *HMLSL-Perfect*

begin

context *hmlsl-perfect*

begin

interpretation *hmlsl* : *hmlsl perfect* :: *cars* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real*

$\langle proof \rangle$

notation $hmlsl.re$ ($\langle re'(-) \rangle$)

notation $hmlsl.cl$ ($\langle cl'(-) \rangle$)

notation $hmlsl.len$ ($\langle len \rangle$)

Safety in the context of HMLSL means the absence of overlapping reservations. Using the somewhere modality, this is easy to formalise.

abbreviation $safe::cars \Rightarrow \sigma$

where $safe\ e \equiv \forall\ c. \neg(c = e) \rightarrow \neg \langle re(c) \wedge re(e) \rangle$

The distance controller ensures, that as long as the cars do not try to change their lane, they keep their distance. More formally, if the reservations of two cars do not overlap, they will also not overlap after an arbitrary amount of time passed. Observe that the cars are allowed to change their dynamical behaviour, i.e., to accelerate and brake.

abbreviation $DC::\sigma$

where $DC \equiv \mathbf{G}(\forall\ c\ d. \neg(c = d) \rightarrow \neg \langle re(c) \wedge re(d) \rangle \rightarrow \Box \tau \neg \langle re(c) \wedge re(d) \rangle)$

To identify possibly dangerous situations during a lane change manoeuvre, we use the *potential collision check*. It allows us to identify situations, where the claim of a car d overlaps with any part of the car c .

abbreviation $pcc::cars \Rightarrow cars \Rightarrow \sigma$

where $pcc\ c\ d \equiv \neg(c = d) \wedge \langle cl(d) \wedge (re(c) \vee cl(c)) \rangle$

The only restriction the lane change controller imposes onto the cars is that in the case of a potential collision, they are not allowed to change the claim into a reservation.

abbreviation $LC::\sigma$

where $LC \equiv \mathbf{G}(\forall\ d. (\exists\ c. pcc\ c\ d) \rightarrow \Box r(d) \perp)$

The safety theorem is as follows. If the controllers of all cars adhere to the specifications given by LC and DC , and we start with an initially safe traffic snapshot, then all reachable traffic snapshots are also safe.

theorem $safety: \models (\forall\ e. safe\ e) \wedge DC \wedge LC \rightarrow \mathbf{G}(\forall\ e. safe\ e)$

$\langle proof \rangle$

While the safety theorem was only proven for a single car, we can show that the choice of this car is irrelevant. That is, if we have a safe situation, and switch the perspective to another car, the resulting situation is also safe.

lemma $safety-switch-invariant: \models (\forall\ e. safe(e)) \rightarrow @c (\forall\ e. safe(e))$

$\langle proof \rangle$

end

end

14 Regular Sensors

This section contains an instantiations of the sensor function for "regular sensors". That is, each car can perceive its own physical size and braking distance. However, it can only perceive the physical size of other cars, and does not know about their braking distance.

```
theory Regular-Sensors
imports ../Length
begin
```

```
definition regular::cars  $\Rightarrow$  traffic  $\Rightarrow$  cars  $\Rightarrow$  real
where regular e ts c  $\equiv$ 
  if (e = c) then traffic.physical-size ts c + traffic.braking-distance ts c
  else traffic.physical-size ts c
```

```
locale regular-sensors = traffic + view
begin
```

```
interpretation regular-sensors: sensors regular :: cars  $\Rightarrow$  traffic  $\Rightarrow$  cars  $\Rightarrow$  real
 $\langle$ proof $\rangle$ 
```

```
notation regular-sensors.space ( $\langle$ space $\rangle$ )
notation regular-sensors.len ( $\langle$ len $\rangle$ )
```

Similar to the situation with perfect sensors, we can show that the perceived length of a car is independent of the spatial transitions between traffic snapshots. The length may only change during evolutions, in particular if the car changes its dynamical behaviour.

```
lemma create-reservation-length-stable:
  (ts - r(d)  $\rightarrow$  ts')  $\rightarrow$  len v ts c = len v ts' c
 $\langle$ proof $\rangle$ 
```

```
lemma create-claim-length-stable:
  (ts - c(d, n)  $\rightarrow$  ts')  $\rightarrow$  len v ts c = len v ts' c
 $\langle$ proof $\rangle$ 
```

```
lemma withdraw-reservation-length-stable:
  (ts - wdr(d, n)  $\rightarrow$  ts')  $\rightarrow$  len v ts c = len v ts' c
 $\langle$ proof $\rangle$ 
```

```
lemma withdraw-claim-length-stable:
  (ts - wdc(d)  $\rightarrow$  ts')  $\rightarrow$  len v ts c = len v ts' c
 $\langle$ proof $\rangle$ 
```

Since the perceived length of cars depends on the owner of the view, we can now prove how this perception changes if we change the perspective of a view.

```
lemma sensors-le: e  $\neq$  c  $\rightarrow$  regular e ts c < regular c ts c
```

$\langle \text{proof} \rangle$

lemma *sensors-leq*: $\text{regular } e \text{ ts } c \leq \text{regular } c \text{ ts } c$
 $\langle \text{proof} \rangle$

lemma *space-eq*: $\text{own } v = \text{own } v' \longrightarrow \text{space ts } v \text{ } c = \text{space ts } v' \text{ } c$
 $\langle \text{proof} \rangle$

lemma *switch-space-le*: $(\text{own } v) \neq c \wedge (v=c > v') \longrightarrow \text{space ts } v \text{ } c < \text{space ts } v' \text{ } c$
 $\langle \text{proof} \rangle$

lemma *switch-space-leq*: $(v=c > v') \longrightarrow \text{space ts } v \text{ } c \leq \text{space ts } v' \text{ } c$
 $\langle \text{proof} \rangle$
end
end

15 HMLSL for Regular Sensors

Within this section, we instantiate HMLSL for cars with regular sensors.

theory *HMLSL-Regular*
imports *../HMLSL Regular-Sensors*
begin

locale *hmlsl-regular* = *regular-sensors* + *restriction*
begin
interpretation *hmlsl* : *hmlsl regular* :: *cars* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real*
 $\langle \text{proof} \rangle$

notation *hmlsl.re* ($\langle \text{re}'(-') \rangle$)
notation *hmlsl.cl* ($\langle \text{cl}'(-') \rangle$)
notation *hmlsl.len* ($\langle \text{len} \rangle$)

The spatial atoms are dependent of the perspective of the view, hence we cannot prove similar lemmas as for perfect sensors.

However, we can still prove lemmas corresponding to the activity and stability rules of the proof system for MSL [5].

Similar to the situation with perfect sensors, needed to instantiate the sensor function, to ensure that the perceived length does not change during spatial transitions.

lemma *backwards-res-act*:
 $(\text{ts} - r(c) \rightarrow \text{ts}') \wedge (\text{ts}', v \models \text{re}(c)) \longrightarrow (\text{ts}, v \models \text{re}(c) \vee \text{cl}(c))$
 $\langle \text{proof} \rangle$

lemma *backwards-res-act-somewhere*:
 $(\text{ts} - r(c) \rightarrow \text{ts}') \wedge (\text{ts}', v \models \langle \text{re}(c) \rangle) \longrightarrow (\text{ts}, v \models \langle \text{re}(c) \vee \text{cl}(c) \rangle)$
 $\langle \text{proof} \rangle$

lemma *backwards-res-stab*:
 $(ts - r(d) \rightarrow ts') \wedge (d \neq c) \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$
 $\langle proof \rangle$

lemma *backwards-c-res-stab*:
 $(ts - c(d, n) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$
 $\langle proof \rangle$

lemma *backwards-wdc-res-stab*:
 $(ts - wdc(d) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$
 $\langle proof \rangle$

lemma *backwards-wdr-res-stab*:
 $(ts - wdr(d, n) \rightarrow ts') \wedge (ts', v \models re(c)) \longrightarrow (ts, v \models re(c))$
 $\langle proof \rangle$

We now proceed to prove the *reservation lemma*, which was crucial in the manual safety proof [2].

lemma *reservation1*: $\models (re(c) \vee cl(c)) \rightarrow \Box r(c) re(c)$
 $\langle proof \rangle$

lemma *reservation2*: $\models (\Box r(c) re(c)) \rightarrow (re(c) \vee cl(c))$
 $\langle proof \rangle$

lemma *reservation*: $\models (\Box r(c) re(c)) \leftrightarrow (re(c) \vee cl(c))$
 $\langle proof \rangle$

end
end

16 Safety for Cars with Regular Sensors

This section contains the definition of requirements for lane change and distance controllers for cars, with the assumption of regular sensors. Using these definitions, we show that safety is an invariant along all possible behaviour of cars. However, we need to slightly amend our notion of safety, compared to the safety proof for perfect sensors.

theory *Safety-Regular*
imports *HMLSL-Regular*
begin
context *hmlsl-regular*
begin

interpretation *hmlsl* : *hmlsl regular* :: *cars* \Rightarrow *traffic* \Rightarrow *cars* \Rightarrow *real*
 $\langle proof \rangle$

notation *hmlsl.space* ($\langle space \rangle$)

notation *hmlsl.re* ($\langle re'(-) \rangle$)

notation $hmlsl.cl(\langle cl'(-) \rangle)$

notation $hmlsl.len(\langle len \rangle)$

First we show that the same "safety" theorem as for perfect sensors can be proven. However, we will subsequently show that this theorem does not ensure safety from the perspective of each car.

The controller definitions for this "flawed" safety are the same as for perfect sensors.

abbreviation $safe::cars \Rightarrow \sigma$

where $safe\ e \equiv \forall\ c. \neg(c = e) \rightarrow \neg\langle re(c) \wedge re(e) \rangle$

abbreviation $DC::\sigma$

where $DC \equiv \mathbf{G}(\forall\ c\ d. \neg(c = d) \rightarrow \neg\langle re(c) \wedge re(d) \rangle \rightarrow \Box\tau \neg\langle re(c) \wedge re(d) \rangle)$

abbreviation $pcc::cars \Rightarrow cars \Rightarrow \sigma$

where $pcc\ c\ d \equiv \neg(c = d) \wedge \langle cl(d) \wedge (re(c) \vee cl(c)) \rangle$

abbreviation $LC::\sigma$

where $LC \equiv \mathbf{G}(\forall\ d. (\exists\ c. pcc\ c\ d) \rightarrow \Box r(d) \perp)$

The safety proof is exactly the same as for perfect sensors. Note in particular, that we fix a single car e for which we show safety.

theorem $safety\text{-}flawed: \models (\forall\ e. safe\ e) \wedge DC \wedge LC \rightarrow \mathbf{G}(\forall\ e. safe\ e)$
 $\langle proof \rangle$

As stated above, the flawed safety theorem does not ensure safety for the perspective of each car. In particular, we can construct a traffic snapshot and a view, such that it satisfies our safety predicate for each car, but if we switch the perspective of the view to another car, the situation is unsafe. A visualisation of this situation can be found in the publication of this work at iFM 2017 [4].

lemma $safety\text{-}not\text{-}invariant\text{-}switch:$

$\exists\ ts\ v. (ts, v \models \forall\ e. safe(e) \wedge (\exists\ c. @c \neg(\forall\ e. safe(e))))$
 $\langle proof \rangle$

Now we show how to amend the controller specifications to gain safety as an invariant even with regular sensors.

The distance controller can be strengthened, by requiring that we switch to the perspective of one of the cars involved first, before checking for the collision. Since all variables are universally quantified, this ensures that no collision exists for the perspective of any car.

abbreviation $DC'::\sigma$

where $DC' \equiv \mathbf{G}(\forall\ c\ d. \neg(c = d) \rightarrow (@d \neg\langle re(c) \wedge re(d) \rangle) \rightarrow \Box\tau @d \neg\langle re(c) \wedge re(d) \rangle)$

The amendment to the lane change controller is slightly different. Instead of checking the potential collision only from the perspective of the car d trying to change lanes, we require that also no other car may perceive a potential collision. Note that the restriction to d 's behaviour can only be enforced within d , if the information from the other car is somehow passed to d . Hence, we require the cars to communicate in some way. However, we do not need to specify, *how* this communication is implemented.

abbreviation $LC'::\sigma$

where $LC' \equiv \mathbf{G} (\forall d. (\exists c. (@c (pcc\ c\ d)) \vee (@d (pcc\ c\ d))) \rightarrow \Box r(d) \perp)$

With these new controllers, we can prove a stronger theorem than before. Instead of proving safety from the perspective of a single car as previously, we now only consider a traffic situation to be safe, if it satisfies the safety predicate from the perspective of *all* cars. Note that this immediately implies the safety invariance theorem proven for perfect sensors.

theorem $safety \models (\forall e. @e (safe\ e)) \wedge DC' \wedge LC' \rightarrow \mathbf{G}(\forall e. @e (safe\ e))$
 $\langle proof \rangle$
end
end

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