

The Hahn and Jordan Decomposition Theorems

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1 Introduction

Signed measures are a generalization of measures that can map measurable sets to negative values. In this work we formalize the Hahn decomposition theorem for signed measures, namely that if $(\Omega, \mathcal{A}, \mu)$ is a measure space for a signed measure μ , then Ω can be decomposed as $\Omega^+ \cup \Omega^-$, where every measurable subset of Ω^+ has a positive measure, and every measurable subset of Ω^- has a negative measure. We then prove that this decomposition is essentially unique, meaning that if $X^+ \cup X^-$ is another such decomposition, then any measurable subset in $(\Omega^+ \Delta X^+) \cup (\Omega^- \Delta X^-)$ has a zero measure.

We also formalize the Jordan decomposition theorem as a corollary, which states that the signed measure μ admits a unique decomposition into a difference $\mu = \mu^+ - \mu^-$ of two positive measures, at least one of which is finite, and such that for any Hahn decomposition $\Omega^+ \cup \Omega^-$ and measurable set A , if $A \subseteq \Omega^-$ then $\mu^+(A) = 0$ and if $A \subseteq \Omega^+$ then $\mu^-(A) = 0$. The formalization is mostly based on [1], Section 16 of Chapter 4.

2 Signed measures

In this section we define signed measures. These are generalizations of measures that can also take negative values but cannot contain both ∞ and $-\infty$ in their range.

2.1 Basic definitions

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theory Hahn-Jordan-Decomposition imports
  HOL-Probability Probability
  Hahn-Jordan-Prelims
begin

definition signed-measure::'a measure ⇒ ('a set ⇒ ereal) ⇒ bool where
  signed-measure M μ ↔ μ {} = 0 ∧ (−∞ ∉ range μ ∨ ∞ ∉ range μ) ∧
  (∀ A. range A ⊆ sets M → disjoint-family A → ∪ (range A) ∈ sets M →
    (λi. μ (A i)) sums μ (∪ (range A))) ∧
  (∀ A. range A ⊆ sets M → disjoint-family A → ∪ (range A) ∈ sets M →
    |μ (∪ (range A))| < ∞ → summable (λi. real-of-ereal |μ (A i)|))
  shows |μ (∪ (range A))| < ∞ → summable (λi. real-of-ereal |μ (A i)|)

lemma signed-measure-empty:
  assumes signed-measure M μ
  shows μ {} = 0 using assms unfolding signed-measure-def by simp

lemma signed-measure-sums:
  assumes signed-measure M μ
  and range A ⊆ M
  and disjoint-family A
  and ∪ (range A) ∈ sets M
  shows (λi. μ (A i)) sums μ (∪ (range A))
  using assms unfolding signed-measure-def by simp

lemma signed-measure-summable:
  assumes signed-measure M μ
  and range A ⊆ M
  and disjoint-family A
  and ∪ (range A) ∈ sets M
  and |μ (∪ (range A))| < ∞
  shows summable (λi. real-of-ereal |μ (A i)|)
  using assms unfolding signed-measure-def by simp

lemma signed-measure-inf-sum:
  assumes signed-measure M μ
  and range A ⊆ M
  and disjoint-family A
  and ∪ (range A) ∈ sets M
  shows (∑ i. μ (A i)) = μ (∪ (range A)) using sums-unique assms
  signed-measure-sums by (metis)

```

```

lemma signed-measure-abs-convergent:
assumes signed-measure M μ
  and range A ⊆ sets M
  and disjoint-family A
  and ∪ (range A) ∈ sets M
  and |μ (∪ (range A))| < ∞
shows summable (λi. real-of-ereal |μ (A i)|) using assms
unfolding signed-measure-def by simp

lemma signed-measure-additive:
assumes signed-measure M μ
shows additive M μ
proof (auto simp add: additive-def)
fix x y
assume x: x ∈ M and y: y ∈ M and x ∩ y = {}
hence disjoint-family (binaryset x y)
  by (auto simp add: disjoint-family-on-def binaryset-def)
have (λi. μ ((binaryset x y) i)) sums (μ x + μ y) using binaryset-sums
  signed-measure-empty[of M μ] assms by simp
have range (binaryset x y) = {x, y, {}} using range-binaryset-eq by simp
moreover have {x, y, {}} ⊆ M using x y by auto
moreover have x ∪ y ∈ sets M using x y by simp
moreover have (∪ (range (binaryset x y))) = x ∪ y
  by (simp add: calculation(1))
ultimately have (λi. μ ((binaryset x y) i)) sums μ (x ∪ y) using assms x y
  signed-measure-empty[of M μ] signed-measure-sums[of M μ]
  ⟨disjoint-family (binaryset x y)⟩ by (metis)
then show μ (x ∪ y) = μ x + μ y
  using ⟨(λi. μ ((binaryset x y) i)) sums (μ x + μ y)⟩ sums-unique2 by force
qed

lemma signed-measure-add:
assumes signed-measure M μ
  and a ∈ sets M
  and b ∈ sets M
  and a ∩ b = {}
shows μ (a ∪ b) = μ a + μ b using additiveD[OF signed-measure-additive]
  assms by auto

lemma signed-measure-disj-sum:
shows finite I ==> signed-measure M μ ==> disjoint-family-on A I ==>
  (λi. i ∈ I ==> A i ∈ sets M) ==> μ (∪ i ∈ I. A i) = (∑ i ∈ I. μ (A i))
proof (induct rule:finite-induct)
case empty
then show ?case unfolding signed-measure-def by simp
next
case (insert x F)
have μ (∪ (A ` insert x F)) = μ ((∪ (A ` F)) ∪ A x)
  by (simp add: Un-commute)

```

```

also have ... =  $\mu(\bigcup(A \setminus F)) + \mu(A \cap x)$ 
proof -
  have  $(\bigcup(A \setminus F)) \cap (A \cap x) = \{\}$  using insert
    by (metis disjoint-family-on-insert inf-commute)
  moreover have  $\bigcup(A \setminus F) \in \text{sets } M$  using insert by auto
  moreover have  $A \cap x \in \text{sets } M$  using insert by simp
  ultimately show ?thesis by (meson insert.preds(1) signed-measure-add)
qed
also have ... =  $(\sum_{i \in F} \mu(A \cap i)) + \mu(A \cap x)$  using insert
  by (metis disjoint-family-on-insert insert-iff)
also have ... =  $(\sum_{i \in \text{insert } x \cap F} \mu(A \cap i))$ 
  by (simp add: add.commute insert.hyps(1) insert.hyps(2))
finally show ?case .
qed

```

```

lemma pos-signed-measure-count-additive:
assumes signed-measure M μ
  and  $\forall E \in \text{sets } M. 0 \leq \mu E$ 
shows countably-additive (sets M) ( $\lambda A. \text{ennreal}(\mu A)$ )
  unfolding countably-additive-def
proof (intro allI impI)
fix A::nat  $\Rightarrow$  'a set
assume range A ⊆ sets M
  and disjoint-family A
  and  $\bigcup(\text{range } A) \in \text{sets } M$  note Aprops = this
have eq:  $\bigwedge i. \mu(A \cap i) = \text{ennreal}(\text{ennreal}(\mu(A \cap i)))$ 
  using assms ennreal-e2ennreal Aprops by simp
have  $(\lambda n. \sum_{i \leq n} \mu(A \cap i)) \longrightarrow \mu(\bigcup(\text{range } A))$  using
  sums-def-le[of λi. μ(A ∩ i) μ(∪(range A))] assms
  signed-measure-sums[of M] Aprops by simp
hence  $((\lambda n. \text{ennreal}(\sum_{i \leq n} \mu(A \cap i))) \longrightarrow$ 
   $\text{ennreal}(\mu(\bigcup(\text{range } A))))$  sequentially
  using tendsto-e2ennrealI[of  $(\lambda n. \sum_{i \leq n} \mu(A \cap i)) \mu(\bigcup(\text{range } A))$ ]
  by simp
moreover have  $\bigwedge n. \text{ennreal}(\sum_{i \leq n} \mu(A \cap i)) = (\sum_{i \leq n} \text{ennreal}(\mu(A \cap i)))$ 
  using e2ennreal-finite-sum by (metis ennreal-nonneg eq finite-atMost)
ultimately have  $((\lambda n. (\sum_{i \leq n} \text{ennreal}(\mu(A \cap i)))) \longrightarrow$ 
   $\text{ennreal}(\mu(\bigcup(\text{range } A))))$  sequentially by simp
hence  $(\lambda i. \text{ennreal}(\mu(A \cap i))) \text{ sums } \text{ennreal}(\mu(\bigcup(\text{range } A)))$ 
  using sums-def-le[of λi. e2ennreal(μ(A ∩ i)) e2ennreal(μ(∪(range A)))]
  by simp
thus  $(\sum i. \text{ennreal}(\mu(A \cap i))) = \text{ennreal}(\mu(\bigcup(\text{range } A)))$ 
  using sums-unique assms by (metis)
qed

```

```

lemma signed-measure-minus:
assumes signed-measure M μ
shows signed-measure M ( $\lambda A. -\mu A$ ) unfolding signed-measure-def

```

```

proof (intro conjI)
show - μ {} = 0 using assms unfolding signed-measure-def by simp
show - ∞ ∉ range (λA. - μ A) ∨ ∞ ∉ range (λA. - μ A)
proof (cases ∞ ∈ range μ)
  case True
    hence -∞ ∉ range μ using assms unfolding signed-measure-def by simp
    hence ∞ ∉ range (λA. - μ A) using ereal-uminus-eq-reorder by blast
    thus - ∞ ∉ range (λA. - μ A) ∨ ∞ ∉ range (λA. - μ A) by simp
next
  case False
    hence -∞ ∉ range (λA. - μ A) using ereal-uminus-eq-reorder
    by (simp add: image-iff)
    thus - ∞ ∉ range (λA. - μ A) ∨ ∞ ∉ range (λA. - μ A) by simp
qed
show ∀ A. range A ⊆ sets M —> disjoint-family A —> ⋃ (range A) ∈ sets M
—>
| - μ (⋃ (range A)) | < ∞ —> summable (λi. real-of-ereal |- μ (A i) |)
proof (intro allI impI)
  fix A::nat ⇒ 'a set
  assume range A ⊆ sets M and disjoint-family A and ⋃ (range A) ∈ sets M
  and | - μ (⋃ (range A)) | < ∞
  thus summable (λi. real-of-ereal |- μ (A i) |) using assms
    unfolding signed-measure-def by simp
qed
show ∀ A. range A ⊆ sets M —> disjoint-family A —> ⋃ (range A) ∈ sets M
—>
(λi. - μ (A i)) sums - μ (⋃ (range A))
proof -
  {
    fix A::nat ⇒ 'a set
    assume range A ⊆ sets M and disjoint-family A and
      ⋃ (range A) ∈ sets M note Aprops = this
    have - ∞ ∉ range (λi. μ (A i)) ∨ ∞ ∉ range (λi. μ (A i))
    proof -
      have range (λi. μ (A i)) ⊆ range μ by auto
      thus ?thesis using assms unfolding signed-measure-def by auto
    qed
    moreover have (λi. μ (A i)) sums μ (⋃ (range A))
      using signed-measure-sums[of M] Aprops assms by simp
    ultimately have (λi. - μ (A i)) sums - μ (⋃ (range A))
      using sums-minus'[of λi. μ (A i)] by simp
  }
  thus ?thesis by auto
qed
qed

locale near-finite-function =
  fixes μ:: 'b set ⇒ ereal
  assumes inf-range: - ∞ ∉ range μ ∨ ∞ ∉ range μ

```

```

lemma (in near-finite-function) finite-subset:
assumes |μ E| < ∞
and A ⊆ E
and μ E = μ A + μ (E - A)
shows |μ A| < ∞
proof (cases ∞ ∈ range μ)
case False
show ?thesis
proof (cases 0 < μ A)
case True
hence |μ A| = μ A by simp
also have ... < ∞ using False by (metis ereal-less-PInfty rangeI)
finally show ?thesis .
next
case False
hence |μ A| = -μ A using not-less-iff-gr-or-eq by fastforce
also have ... = μ (E - A) - μ E
proof -
have μ E = μ A + μ (E - A) using assms by simp
hence μ E - μ A = μ (E - A)
by (metis abs-ereal-uminus assms(1) calculation ereal-diff-add-inverse
ereal-infty-less(2) ereal-minus(5) ereal-minus-less-iff
ereal-minus-less-minus ereal-uminus-uminus less-ereal.simps(2)
minus-ereal-def plus-ereal.simps(3))
thus ?thesis using assms(1) ereal-add-uminus-conv-diff ereal-eq-minus
by auto
qed
also have ... ≤ μ (E - A) + |μ E|
by (metis ‹- μ A = μ (E - A) - μ E› abs-ereal-less0 abs-ereal-pos
ereal-diff-le-self add-increasing2 less-eq-ereal-def
minus-ereal-def not-le-imp-less)
also have ... < ∞ using assms ‹∞ ∉ range μ›
by (metis UNIV-I ereal-less-PInfty ereal-plus-eq-PInfty image-eqI)
finally show ?thesis .
qed
next
case True
hence -∞ ∉ range μ using inf-range by simp
hence -∞ < μ A by (metis ereal-infty-less(2) rangeI)
show ?thesis
proof (cases μ A < 0)
case True
hence |μ A| = -μ A using not-less-iff-gr-or-eq by fastforce
also have ... < ∞ using ‹-∞ < μ A› using ereal-uminus-less-reorder
by blast
finally show ?thesis .
next
case False

```

```

hence  $|\mu A| = \mu A$  by simp
also have ... =  $\mu E - \mu (E - A)$ 
proof -
  have  $\mu E = \mu A + \mu (E - A)$  using assms by simp
  thus  $\mu A = \mu E - \mu (E - A)$  by (metis add.right-neutral assms(1)
    add-diff-eq-ereal calculation ereal-diff-add-eq-diff-diff-swap
    ereal-diff-add-inverse ereal-infity-less(1) ereal-plus-eq-PInfty
    ereal-x-minus-x)
qed
also have ...  $\leq |\mu E| - \mu (E - A)$ 
by (metis <|μ A| = μ A> <μ A = μ E - μ (E - A)> abs-ereal-ge0
  abs-ereal-pos abs-ereal-uminus antisym-conv ereal-0-le-uminus-iff
  ereal-abs-diff ereal-diff-le-mono-left ereal-diff-le-self le-cases
  less-eq-ereal-def minus-ereal-def)
also have ...  $< \infty$ 
proof -
  have  $-\infty < \mu (E - A)$  using < $-\infty \notin range \mu$ >
  by (metis ereal-infity-less(2) rangeI)
  hence  $-\mu (E - A) < \infty$  using ereal-uminus-less-reorder by blast
  thus ?thesis using assms by (simp add: ereal-minus-eq-PInfty-iff
    ereal-uminus-eq-reorder)
qed
finally show ?thesis .
qed
qed

locale signed-measure-space=
  fixes M::'a measure and μ
  assumes sgn-meas: signed-measure M μ

sublocale signed-measure-space ⊆ near-finite-function
proof (unfold-locales)
  show  $-\infty \notin range \mu \vee \infty \notin range \mu$  using sgn-meas
    unfolding signed-measure-def by simp
qed

context signed-measure-space
begin
lemma signed-measure-finite-subset:
  assumes E ∈ sets M
  and |μ E| < ∞
  and A ∈ sets M
  and A ⊆ E
  shows |μ A| < ∞
proof (rule finite-subset)
  show |μ E| < ∞ A ⊆ E using assms by auto
  show  $\mu E = \mu A + \mu (E - A)$  using assms
    sgn-meas signed-measure-add[of M μ A E - A]
    by (metis Diff-disjoint Diff-partition sets.Diff)

```

qed

lemma *measure-space-e2ennreal* :

assumes *measure-space (space M) (sets M) m* \wedge $(\forall E \in \text{sets } M. m E < \infty) \wedge (\forall E \in \text{sets } M. m E \geq 0)$

shows $\forall E \in \text{sets } M. e2ennreal (m E) < \infty$

proof

fix *E*

assume *E ∈ sets M*

show *e2ennreal (m E) < ∞*

proof –

have *m E < ∞* **using** *assms {E ∈ sets M}*

by *blast*

then have *e2ennreal (m E) < ∞* **using** *e2ennreal-less-top*

using *{m E < ∞}* **by** *auto*

thus *?thesis by simp*

qed

qed

2.2 Positive and negative subsets

The Hahn decomposition theorem is based on the notions of positive and negative measurable sets. A measurable set is positive (resp. negative) if all its measurable subsets have a positive (resp. negative) measure by μ . The decomposition theorem states that any measure space for a signed measure can be decomposed into a positive and a negative measurable set.

definition *pos-meas-set where*

pos-meas-set E ↔ E ∈ sets M ∧ (forall A ∈ sets M. A ⊆ E → 0 ≤ μ A)

definition *neg-meas-set where*

neg-meas-set E ↔ E ∈ sets M ∧ (forall A ∈ sets M. A ⊆ E → μ A ≤ 0)

lemma *pos-meas-setI*:

assumes *E ∈ sets M*

and $\bigwedge A. A \in \text{sets } M \implies A \subseteq E \implies 0 \leq \mu A$

shows *pos-meas-set E unfolding pos-meas-set-def using assms by simp*

lemma *pos-meas-setD1* :

assumes *pos-meas-set E*

shows *E ∈ sets M*

using *assms unfolding pos-meas-set-def*

by *simp*

lemma *neg-meas-setD1* :

assumes *neg-meas-set E*

shows *E ∈ sets M using assms unfolding neg-meas-set-def by simp*

lemma *neg-meas-setI*:

```

assumes  $E \in \text{sets } M$ 
and  $\bigwedge A. A \in \text{sets } M \implies A \subseteq E \implies \mu A \leq 0$ 
shows neg-meas-set  $E$  unfolding neg-meas-set-def using assms by simp

lemma pos-meas-self:
assumes pos-meas-set  $E$ 
shows  $0 \leq \mu E$  using assms unfolding pos-meas-set-def by simp

lemma empty-pos-meas-set:
shows pos-meas-set  $\{\}$ 
by (metis bot.extremum-uniqueI eq-iff pos-meas-set-def sets.empty-sets
sgn-meas signed-measure-empty)

lemma empty-neg-meas-set:
shows neg-meas-set  $\{\}$ 
by (metis neg-meas-set-def order-refl sets.empty-sets sgn-meas
signed-measure-empty subset-empty)

lemma pos-measure-meas:
assumes pos-meas-set  $E$ 
and  $A \subseteq E$ 
and  $A \in \text{sets } M$ 
shows  $0 \leq \mu A$  using assms unfolding pos-meas-set-def by simp

lemma pos-meas-subset:
assumes pos-meas-set  $A$ 
and  $B \subseteq A$ 
and  $B \in \text{sets } M$ 
shows pos-meas-set  $B$  using assms pos-meas-set-def by auto

lemma neg-meas-subset:
assumes neg-meas-set  $A$ 
and  $B \subseteq A$ 
and  $B \in \text{sets } M$ 
shows neg-meas-set  $B$  using assms neg-meas-set-def by auto

lemma pos-meas-set-Union:
assumes  $\bigwedge (i:\text{nat}). \text{pos-meas-set} (A i)$ 
and  $\bigwedge i. A i \in \text{sets } M$ 
and  $|\mu (\bigcup i. A i)| < \infty$ 
shows pos-meas-set  $(\bigcup i. A i)$ 
proof (rule pos-meas-setI)
show  $\bigcup (\text{range } A) \in \text{sets } M$  using sigma-algebra.countable-UN assms by simp
obtain  $B$  where disjoint-family  $B$  and  $(\bigcup (i:\text{nat}). B i) = (\bigcup (i:\text{nat}). A i)$ 
and  $\bigwedge i. B i \in \text{sets } M$  and  $\bigwedge i. B i \subseteq A i$  using disj-Union2 assms by auto
fix  $C$ 
assume  $C \in \text{sets } M$  and  $C \subseteq (\bigcup i. A i)$ 
hence  $C = C \cap (\bigcup i. A i)$  by auto
also have ... =  $C \cap (\bigcup i. B i)$  using  $\langle (\bigcup i. B i) = (\bigcup i. A i) \rangle$  by simp

```

```

also have ... = ( $\bigcup i. C \cap B_i$ ) by auto
finally have  $C = (\bigcup i. C \cap B_i)$  .
hence  $\mu C = \mu (\bigcup i. C \cap B_i)$  by simp
also have ... = ( $\sum i. \mu (C \cap (B_i))$ )
proof (rule signed-measure-inf-sum[symmetric])
  show signed-measure  $M \mu$  using sgn-meas by simp
  show disjoint-family ( $\lambda i. C \cap B_i$ ) using disjoint-family B
    by (meson Int-iff disjoint-family-subset subset-iff)
  show range ( $\lambda i. C \cap B_i$ )  $\subseteq$  sets  $M$  using  $\langle C \in \text{sets } M \rangle \langle \bigwedge i. B_i \in \text{sets } M \rangle$ 
    by auto
  show  $(\bigcup i. C \cap B_i) \in \text{sets } M$  using  $\langle C = (\bigcup i. C \cap B_i) \rangle \langle C \in \text{sets } M \rangle$ 
    by simp
qed
also have ...  $\geq 0$ 
proof (rule suminf-nonneg)
  show  $\bigwedge n. 0 \leq \mu (C \cap B_n)$ 
proof -
  fix  $n$ 
  have  $C \cap B_n \subseteq A_n$  using  $\langle \bigwedge i. B_i \subseteq A_i \rangle$  by auto
  moreover have  $C \cap B_n \in \text{sets } M$  using  $\langle C \in \text{sets } M \rangle \langle \bigwedge i. B_i \in \text{sets } M \rangle$ 
    by simp
  ultimately show  $0 \leq \mu (C \cap B_n)$  using assms pos-measure-meas[of  $A_n$ ]
    by simp
qed
have summable ( $\lambda i. \text{real-of-ereal} (\mu (C \cap B_i))$ )
proof (rule summable-norm-cancel)
  have  $\bigwedge n. \text{norm} (\text{real-of-ereal} (\mu (C \cap B_n))) =$ 
    real-of-ereal  $|\mu (C \cap B_n)|$  by simp
  moreover have summable ( $\lambda i. \text{real-of-ereal} |\mu (C \cap B_i)|$ )
  proof (rule signed-measure-abs-convergent)
    show signed-measure  $M \mu$  using sgn-meas by simp
    show range ( $\lambda i. C \cap B_i$ )  $\subseteq$  sets  $M$  using  $\langle C \in \text{sets } M \rangle \langle \bigwedge i. B_i \in \text{sets } M \rangle$ 
    show disjoint-family ( $\lambda i. C \cap B_i$ ) using disjoint-family B
      by (meson Int-iff disjoint-family-subset subset-iff)
    show  $(\bigcup i. C \cap B_i) \in \text{sets } M$  using  $\langle C = (\bigcup i. C \cap B_i) \rangle \langle C \in \text{sets } M \rangle$ 
      by simp
    have  $|\mu C| < \infty$ 
    proof (rule signed-measure-finite-subset)
      show  $(\bigcup i. A_i) \in \text{sets } M$  using assms by simp
      show  $|\mu (\bigcup (range A))| < \infty$  using assms by simp
      show  $C \in \text{sets } M$  using  $\langle C \in \text{sets } M \rangle$  .
      show  $C \subseteq \bigcup (range A)$  using  $\langle C \subseteq \bigcup (range A) \rangle$  .
    qed
    thus  $|\mu (\bigcup i. C \cap B_i)| < \infty$  using  $\langle C = (\bigcup i. C \cap B_i) \rangle$  by simp
  qed
  ultimately show summable ( $\lambda n. \text{norm} (\text{real-of-ereal} (\mu (C \cap B_n)))$ )
    by auto
qed

```

```

thus summable ( $\lambda i. \mu (C \cap B i)$ ) by (simp add:  $\langle \bigwedge n. 0 \leq \mu (C \cap B n) \rangle$ 
summable-ereal-pos)
qed
finally show  $0 \leq \mu C$  .
qed

lemma pos-meas-set-pos-lim:
assumes  $\bigwedge (i::nat). \text{pos-meas-set} (A i)$ 
and  $\bigwedge i. A i \in \text{sets } M$ 
shows  $0 \leq \mu (\bigcup i. A i)$ 
proof -
obtain B where disjoint-family B and  $(\bigcup (i::nat). B i) = (\bigcup (i::nat). A i)$ 
and  $\bigwedge i. B i \in \text{sets } M$  and  $\bigwedge i. B i \subseteq A i$  using disj-Union2 assms by auto
note Bprops = this
have sums:  $(\lambda n. \mu (B n))$  sums  $\mu (\bigcup i. B i)$ 
proof (rule signed-measure-sums)
show signed-measure M  $\mu$  using sgn-meas .
show range B  $\subseteq$  sets M using Bprops by auto
show disjoint-family B using Bprops by simp
show  $\bigcup (\text{range } B) \in \text{sets } M$  using Bprops by blast
qed
hence summable ( $\lambda n. \mu (B n)$ ) using sumssummable[of  $\lambda n. \mu (B n)$ ] by simp
hence suminf ( $\lambda n. \mu (B n)$ ) =  $\mu (\bigcup i. B i)$  using sums sums-iff by auto
thus ?thesis using suminf-nonneg
by (metis Bprops(2) Bprops(3) Bprops(4) ⟨summable ( $\lambda n. \mu (B n)$ )⟩ assms(1))

pos-measure-meas)
qed

lemma pos-meas-disj-union:
assumes pos-meas-set A
and pos-meas-set B
and  $A \cap B = \{\}$ 
shows pos-meas-set ( $A \cup B$ ) unfolding pos-meas-set-def
proof (intro conjI ballI impI)
show  $A \cup B \in \text{sets } M$ 
by (metis assms(1) assms(2) pos-meas-set-def sets.Un)
next
fix C
assume  $C \in \text{sets } M$  and  $C \subseteq A \cup B$ 
define DA where  $DA = C \cap A$ 
define DB where  $DB = C \cap B$ 
have DA ∈ sets M using DA-def ⟨C ∈ sets M⟩ assms(1) pos-meas-set-def
by blast
have DB ∈ sets M using DB-def ⟨C ∈ sets M⟩ assms(2) pos-meas-set-def
by blast
have  $DA \cap DB = \{\}$  unfolding DA-def DB-def using assms by auto
have  $C = DA \cup DB$  unfolding DA-def DB-def using ⟨C ⊆ A ∪ B⟩ by auto
have  $0 \leq \mu DB$  using assms unfolding DB-def pos-meas-set-def

```

```

    by (metis DB-def Int-lower2 ‹DB ∈ sets M›)
also have ... ≤ μ DA + μ DB using assms unfolding pos-meas-set-def
    by (metis DA-def Diff-Diff-Int Diff-subset Int-commute ‹DA ∈ sets M›
         ereal-le-add-self2)
also have ... = μ C using signed-measure-add sgn-meas ‹DA ∈ sets M›
    ‹DB ∈ sets M› ‹DA ∩ DB = {}› ‹C = DA ∪ DB› by metis
finally show 0 ≤ μ C .
qed

```

```

lemma pos-meas-set-union:
assumes pos-meas-set A
and pos-meas-set B
shows pos-meas-set (A ∪ B)
proof -
define C where C = B - A
have A ∪ C = A ∪ B unfolding C-def by auto
moreover have pos-meas-set (A ∪ C)
proof (rule pos-meas-disj-union)
show pos-meas-set C unfolding C-def
by (meson Diff-subset assms(1) assms(2) sets.Diff
signed-measure-space.pos-meas-set-def
signed-measure-space.pos-meas-subset signed-measure-space-axioms)
show pos-meas-set A using assms by simp
show A ∩ C = {} unfolding C-def by auto
qed
ultimately show ?thesis by simp
qed

```

```

lemma neg-meas-disj-union:
assumes neg-meas-set A
and neg-meas-set B
and A ∩ B = {}
shows neg-meas-set (A ∪ B) unfolding neg-meas-set-def
proof (intro conjI ballI impI)
show A ∪ B ∈ sets M
by (metis assms(1) assms(2) neg-meas-set-def sets.Un)
next
fix C
assume C ∈ sets M and C ⊆ A ∪ B
define DA where DA = C ∩ A
define DB where DB = C ∩ B
have DA ∈ sets M using DA-def ‹C ∈ sets M› assms(1) neg-meas-set-def
by blast
have DB ∈ sets M using DB-def ‹C ∈ sets M› assms(2) neg-meas-set-def
by blast
have DA ∩ DB = {} unfolding DA-def DB-def using assms by auto
have C = DA ∪ DB unfolding DA-def DB-def using ‹C ⊆ A ∪ B› by auto
have μ C = μ DA + μ DB using signed-measure-add sgn-meas ‹DA ∈ sets M›
    ‹DB ∈ sets M› ‹DA ∩ DB = {}› ‹C = DA ∪ DB› by metis

```

```

also have ... ≤ μ DB using assms unfolding neg-meas-set-def
  by (metis DA-def Int-lower2 ⟨DA ∈ sets M⟩ add-decreasing dual-order.refl)
also have ... ≤ 0 using assms unfolding DB-def neg-meas-set-def
  by (metis DB-def Int-lower2 ⟨DB ∈ sets M⟩)
finally show μ C ≤ 0 .
qed

```

```

lemma neg-meas-set-union:
  assumes neg-meas-set A
    and neg-meas-set B
  shows neg-meas-set (A ∪ B)
proof -
  define C where C = B - A
  have A ∪ C = A ∪ B unfolding C-def by auto
  moreover have neg-meas-set (A ∪ C)
  proof (rule neg-meas-disj-union)
    show neg-meas-set C unfolding C-def
      by (meson Diff-subset assms(1) assms(2) sets.Diff neg-meas-set-def
           neg-meas-subset signed-measure-space-axioms)
    show neg-meas-set A using assms by simp
    show A ∩ C = {} unfolding C-def by auto
  qed
  ultimately show ?thesis by simp
qed

```

```

lemma neg-meas-self :
  assumes neg-meas-set E
  shows μ E ≤ 0 using assms unfolding neg-meas-set-def by simp

```

```

lemma pos-meas-set-opp:
  assumes signed-measure-space.pos-meas-set M (λ A. - μ A) A
  shows neg-meas-set A
proof -
  have m-meas-pos : signed-measure M (λ A. - μ A)
    using assms signed-measure-space-def
    by (simp add: sgn-meas signed-measure-minus)
  thus ?thesis
    by (metis assms ereal-0-le-uminus-iff neg-meas-setI
          signed-measure-space.intro signed-measure-space.pos-meas-set-def)
qed

```

```

lemma neg-meas-set-opp:
  assumes signed-measure-space.neg-meas-set M (λ A. - μ A) A
  shows pos-meas-set A
proof -
  have m-meas-neg : signed-measure M (λ A. - μ A)
    using assms signed-measure-space-def
    by (simp add: sgn-meas signed-measure-minus)
  thus ?thesis

```

```

by (metis assms ereal-uminus-le-0-iff m-meas-neg pos-meas-setI
      signed-measure-space.intro signed-measure-space.neg-meas-set-def)
qed
end

lemma signed-measure-inter:
assumes signed-measure M μ
and A ∈ sets M
shows signed-measure M (λE. μ (E ∩ A)) unfolding signed-measure-def
proof (intro conjI)
show μ ({}) ∩ A = 0 using assms(1) signed-measure-empty by auto
show -∞ ∉ range (λE. μ (E ∩ A)) ∨ ∞ ∉ range (λE. μ (E ∩ A))
proof (rule ccontr)
assume ¬ (-∞ ∉ range (λE. μ (E ∩ A)) ∨ ∞ ∉ range (λE. μ (E ∩ A)))
hence -∞ ∈ range (λE. μ (E ∩ A)) ∧ ∞ ∈ range (λE. μ (E ∩ A)) by simp
hence -∞ ∈ range μ ∧ ∞ ∈ range μ by auto
thus False using assms unfolding signed-measure-def by simp
qed
show ∀ E. range E ⊆ sets M → disjoint-family E → ∪ (range E) ∈ sets M
→
  (λi. μ (E i ∩ A)) sums μ (∪ (range E) ∩ A)
proof (intro allI impI)
fix E::nat ⇒ 'a set
assume range E ⊆ sets M and disjoint-family E and ∪ (range E) ∈ sets M
note Eprops = this
define F where F = (λi. E i ∩ A)
have (λi. μ (F i)) sums μ (∪ (range F))
proof (rule signed-measure-sums)
show signed-measure M μ using assms by simp
show range F ⊆ sets M using Eprops F-def assms by blast
show disjoint-family F using Eprops F-def assms
by (metis disjoint-family-subset inf.absorb-iff2 inf-commute
     inf-right-idem)
show ∪ (range F) ∈ sets M using Eprops assms unfolding F-def
  by (simp add: Eprops assms countable-Un-Int(1) sets.Int)
qed
moreover have ∪ (range F) = A ∩ ∪ (range E) unfolding F-def by auto
ultimately show (λi. μ (E i ∩ A)) sums μ (∪ (range E) ∩ A)
  unfolding F-def by simp
qed
show ∀ E. range E ⊆ sets M →
disjoint-family E →
  ∪ (range E) ∈ sets M → |μ (∪ (range E) ∩ A)| < ∞ →
    summable (λi. real-of-ereal |μ (E i ∩ A)|)
proof (intro allI impI)
fix E::nat ⇒ 'a set
assume range E ⊆ sets M and disjoint-family E and
  ∪ (range E) ∈ sets M and |μ (∪ (range E) ∩ A)| < ∞ note Eprops = this
show summable (λi. real-of-ereal |μ (E i ∩ A)|)

```

```

proof (rule signed-measure-summable)
  show signed-measure M μ using assms by simp
  show range (λi. E i ∩ A) ⊆ sets M using Eprops assms by blast
  show disjoint-family (λi. E i ∩ A) using Eprops assms
    disjoint-family-subset inf.absorb-iff2 inf-commute inf-right-idem
    by fastforce
  show (∪ i. E i ∩ A) ∈ sets M using Eprops assms
    by (simp add: Eprops assms countable-Un-Int(1) sets.Int)
  show |μ (∪ i. E i ∩ A)| < ∞ using Eprops by auto
  qed
  qed
qed

context signed-measure-space
begin

lemma pos-signed-to-meas-space :
  assumes pos-meas-set M1
  and m1 = (λA. μ (A ∩ M1))
  shows measure-space (space M) (sets M) m1 unfolding measure-space-def
  proof (intro conjI)
    show sigma-algebra (space M) (sets M)
    by (simp add: sets.sigma-algebra-axioms)
    show positive (sets M) m1 using assms unfolding pos-meas-set-def
    by (metis Sigma-Algebra.positive-def Un-Int-eq(4)
      e2ennreal-neg neg-meas-self sup-bot-right empty-neg-meas-set)
    show countably-additive (sets M) m1
    proof (rule pos-signed-measure-count-additive)
      show ∀ E∈sets M. 0 ≤ m1 E by (metis assms inf.cobounded2
        pos-meas-set-def sets.Int)
      show signed-measure M m1 using assms pos-meas-set-def
        signed-measure-inter[of M μ M1] sgn-meas by blast
    qed
  qed

lemma neg-signed-to-meas-space :
  assumes neg-meas-set M2
  and m2 = (λA. -μ (A ∩ M2))
  shows measure-space (space M) (sets M) m2 unfolding measure-space-def
  proof (intro conjI)
    show sigma-algebra (space M) (sets M)
    by (simp add: sets.sigma-algebra-axioms)
    show positive (sets M) m2 using assms unfolding neg-meas-set-def
    by (metis Sigma-Algebra.positive-def e2ennreal-neg ereal-uminus-zero
      inf.absorb-iff2 inf.orderE inf-bot-right neg-meas-self pos-meas-self
      empty-neg-meas-set empty-pos-meas-set)
    show countably-additive (sets M) m2
    proof (rule pos-signed-measure-count-additive)
      show ∀ E∈sets M. 0 ≤ m2 E
      by (metis assms ereal-uminus-eq-reorder ereal-uminus-le-0-iff)

```

```

inf.cobounded2 neg-meas-set-def sets.Int)
have signed-measure M (λA. μ (A ∩ M2)) using assms neg-meas-set-def
  signed-measure-inter[of M μ M2] sgn-meas by blast
  thus signed-measure M m2 using signed-measure-minus assms by simp
qed
qed

lemma pos-part-meas-nul-neg-set :
assumes pos-meas-set M1
  and neg-meas-set M2
  and m1 = (λA. μ (A ∩ M1))
  and E ∈ sets M
  and E ⊆ M2
shows m1 E = 0
proof -
  have m1 E ≥ 0 using assms unfolding pos-meas-set-def
    by (simp add: ‹E ∈ sets M› sets.Int)
  have μ E ≤ 0 using ‹E ⊆ M2› assms unfolding neg-meas-set-def
    using ‹E ∈ sets M› by blast
  then have m1 E ≤ 0 using ‹μ E ≤ 0› assms
    by (metis Int-Un-eq(1) Un-subset-iff ‹E ∈ sets M› ‹E ⊆ M2› pos-meas-setD1
      sets.Int signed-measure-space.neg-meas-set-def
      signed-measure-space-axioms)
  thus m1 E = 0 using ‹m1 E ≥ 0› ‹m1 E ≤ 0› by auto
qed

lemma neg-part-meas-nul-pos-set :
assumes pos-meas-set M1
  and neg-meas-set M2
  and m2 = (λA. -μ (A ∩ M2))
  and E ∈ sets M
  and E ⊆ M1
shows m2 E = 0
proof -
  have m2 E ≥ 0 using assms unfolding neg-meas-set-def
    by (simp add: ‹E ∈ sets M› sets.Int)
  have μ E ≥ 0 using assms unfolding pos-meas-set-def by blast
  then have m2 E ≤ 0 using ‹μ E ≥ 0› assms
    by (metis ‹E ∈ sets M› ‹E ⊆ M1› ereal-0-le-uminus-iff ereal-uminus-uminus
      inf-sup-ord(1) neg-meas-setD1 pos-meas-set-def pos-meas-subset
      sets.Int)
  thus m2 E = 0 using ‹m2 E ≥ 0› ‹m2 E ≤ 0› by auto
qed

definition pos-sets where
pos-sets = {A. A ∈ sets M ∧ pos-meas-set A}

definition pos-img where
pos-img = {μ A | A ∈ pos-sets}

```

2.3 Essential uniqueness

In this part, under the assumption that a measure space for a signed measure admits a decomposition into a positive and a negative set, we prove that this decomposition is essentially unique; in other words, that if two such decompositions (P, N) and (X, Y) exist, then any measurable subset of $(P \Delta X) \cup (N \Delta Y)$ has a null measure.

definition hahn-space-decomp where

hahn-space-decomp M1 M2 \equiv (pos-meas-set M1) \wedge (neg-meas-set M2) \wedge (space M = M1 \cup M2) \wedge (M1 \cap M2 = {})

lemma pos-neg-null-set:

assumes pos-meas-set A
and neg-meas-set A

shows $\mu A = 0$ **using** assms pos-meas-self[of A] neg-meas-self[of A] **by** simp

lemma pos-diff-neg-meas-set:

assumes (pos-meas-set M1)
and (neg-meas-set N2)
and (space M = N1 \cup N2)
and N1 \in sets M

shows neg-meas-set $((M1 - N1) \cap \text{space } M)$ **using** assms neg-meas-subset
by (metis Diff-subset-conv Int-lower2 pos-meas-setD1 sets.Diff

sets.Int-space-eq2)

lemma neg-diff-pos-meas-set:

assumes (neg-meas-set M2)
and (pos-meas-set N1)
and (space M = N1 \cup N2)
and N2 \in sets M

shows pos-meas-set $((M2 - N2) \cap \text{space } M)$

proof –

have $(M2 - N2) \cap \text{space } M \subseteq N1$ **using** assms **by** auto

thus ?thesis **using** assms pos-meas-subset neg-meas-setD1 **by** blast

qed

lemma pos-sym-diff-neg-meas-set:

assumes hahn-space-decomp M1 M2
and hahn-space-decomp N1 N2
shows neg-meas-set $((\text{sym-diff } M1 N1) \cap \text{space } M)$ **using** assms
unfolding hahn-space-decomp-def
by (metis Int-Un-distrib2 neg-meas-set-union pos-meas-setD1
pos-diff-neg-meas-set)

lemma neg-sym-diff-pos-meas-set:

assumes hahn-space-decomp M1 M2
and hahn-space-decomp N1 N2
shows pos-meas-set $((\text{sym-diff } M2 N2) \cap \text{space } M)$ **using** assms
neg-diff-pos-meas-set **unfolding** hahn-space-decomp-def

```

by (metis (no-types, lifting) Int-Un-distrib2 neg-meas-setD1
    pos-meas-set-union)

lemma pos-meas-set-diff:
assumes pos-meas-set A
and B ∈ sets M
shows pos-meas-set ((A − B) ∩ (space M)) using pos-meas-subset
by (metis Diff-subset assms(1) assms(2) pos-meas-setD1 sets.Diff
sets.Int-space-eq2)

lemma pos-meas-set-sym-diff:
assumes pos-meas-set A
and pos-meas-set B
shows pos-meas-set ((sym-diff A B) ∩ space M) using pos-meas-set-diff
by (metis Int-Un-distrib2 assms(1) assms(2) pos-meas-setD1
    pos-meas-set-union)

lemma neg-meas-set-diff:
assumes neg-meas-set A
and B ∈ sets M
shows neg-meas-set ((A − B) ∩ (space M)) using neg-meas-subset
by (metis Diff-subset assms(1) assms(2) neg-meas-setD1 sets.Diff
sets.Int-space-eq2)

lemma neg-meas-set-sym-diff:
assumes neg-meas-set A
and neg-meas-set B
shows neg-meas-set ((sym-diff A B) ∩ space M) using neg-meas-set-diff
by (metis Int-Un-distrib2 assms(1) assms(2) neg-meas-setD1
    neg-meas-set-union)

lemma hahn-decomp-space-diff:
assumes hahn-space-decomp M1 M2
and hahn-space-decomp N1 N2
shows pos-meas-set ((sym-diff M1 N1 ∪ sym-diff M2 N2) ∩ space M)
neg-meas-set ((sym-diff M1 N1 ∪ sym-diff M2 N2) ∩ space M)
proof –
show pos-meas-set ((sym-diff M1 N1 ∪ sym-diff M2 N2) ∩ space M)
by (metis Int-Un-distrib2 assms(1) assms(2) hahn-space-decomp-def
    neg-sym-diff-pos-meas-set pos-meas-set-sym-diff pos-meas-set-union)
show neg-meas-set ((sym-diff M1 N1 ∪ sym-diff M2 N2) ∩ space M)
by (metis Int-Un-distrib2 assms(1) assms(2) hahn-space-decomp-def
    neg-meas-set-sym-diff neg-meas-set-union pos-sym-diff-neg-meas-set)
qed

lemma hahn-decomp-ess-unique:
assumes hahn-space-decomp M1 M2
and hahn-space-decomp N1 N2
and C ⊆ sym-diff M1 N1 ∪ sym-diff M2 N2

```

```

and  $C \in \text{sets } M$ 
shows  $\mu C = 0$ 
proof -
have  $C \subseteq (\text{sym-diff } M1 N1 \cup \text{sym-diff } M2 N2) \cap \text{space } M$  using assms
  by (simp add: sets.sets-into-space)
thus ?thesis using assms hahn-decomp-space-diff pos-neg-null-set
  by (meson neg-meas-subset pos-meas-subset)
qed

```

3 Existence of a positive subset

The goal of this part is to prove that any measurable set of finite and positive measure must contain a positive subset with a strictly positive measure.

3.1 A sequence of negative subsets

definition inf-neg where

```

inf-neg  $A = (\text{if } (A \notin \text{sets } M \vee \text{pos-meas-set } A) \text{ then } (0::nat)$ 
       $\text{else Inf } \{n|n. (1::nat) \leq n \wedge (\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < \text{ereal}(-1/n))\})$ 

```

lemma inf-neg-ne:

assumes $A \in \text{sets } M$

and $\neg \text{pos-meas-set } A$

shows $\{n|n. (1::nat) \leq n \wedge (\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < \text{ereal}(-1/n))\} \neq \{\}$

proof -

```

define  $N$  where  $N = \{n|n. (1::nat) \leq n \wedge$ 
       $(\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < \text{ereal}(-1/n))\}$ 

```

```

have  $\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < 0$  using assms unfolding pos-meas-set-def
  by auto

```

from this obtain B where $B \in \text{sets } M$ and $B \subseteq A$ and $\mu B < 0$ by auto

hence $\exists n|n. (1::nat) \leq n \wedge \mu B < \text{ereal}(-1/n)$

proof (cases $\mu B = -\infty$)

case True

hence $\mu B < -1/(2::nat)$ by simp

thus ?thesis using numeral-le-real-of-nat-iff one-le-numeral by blast

next

case False

hence real-of-ereal ($\mu B < 0$) using $\langle \mu B < 0 \rangle$

```

by (metis Infty-neq-0(3) ereal-mult-eq-MInfty ereal-zero-mult
    less-eq-ereal-def less-eq-real-def less-ereal.simps(2)
    real-of-ereal-eq-0 real-of-ereal-le-0)

```

hence $\exists n|n. Suc 0 \leq n \wedge \text{real-of-ereal } (\mu B) < -1/n$

proof -

```

define nw where  $nw = Suc (\text{nat} (\text{floor} (-1 / (\text{real-of-ereal } (\mu B)))))$ 

```

have $Suc 0 \leq nw$ unfolding nw-def by simp

```

have  $0 < -1 / (\text{real-of-ereal } (\mu B))$  using  $\langle \text{real-of-ereal } (\mu B) < 0 \rangle$ 
  by simp

```

```

have  $-1 / (\text{real-of-ereal } (\mu B)) < nw$  unfolding  $nw\text{-def}$  by linarith
hence  $1/nw < 1/(-1 / (\text{real-of-ereal } (\mu B)))$ 
using  $\langle 0 < -1 / (\text{real-of-ereal } (\mu B)) \rangle$  by (metis frac-less2
 $\text{le-eq-less-or-eq of-nat-1 of-nat-le-iff zero-less-one}$ )
also have ... =  $- (\text{real-of-ereal } (\mu B))$  by simp
finally have  $1/nw < - (\text{real-of-ereal } (\mu B))$  .
hence  $\text{real-of-ereal } (\mu B) < -1/nw$  by simp
thus ?thesis using  $\langle \text{Suc } 0 \leq nw \rangle$  by auto
qed
from this obtain n1::nat where  $\text{Suc } 0 \leq n1$ 
and  $\text{real-of-ereal } (\mu B) < -1/n1$  by auto
hence  $\text{ereal } (\text{real-of-ereal } (\mu B)) < -1/n1$  using real-ereal-leq[of  $\mu B$ ]
 $\langle \mu B < 0 \rangle$  by simp
moreover have  $\mu B = \text{real-of-ereal } (\mu B)$  using  $\langle \mu B < 0 \rangle$  False
by (metis less-ereal.simps(2) real-of-ereal.elims zero-ereal-def)
ultimately show ?thesis using  $\langle \text{Suc } 0 \leq n1 \rangle$  by auto
qed
from this obtain n0::nat where  $(1:\text{nat}) \leq n0$  and  $\mu B < -1/n0$  by auto
hence  $n0 \in \{n:\text{nat} | n. (1:\text{nat}) \leq n \wedge$ 
 $(\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < \text{ereal}(-1/n))\}$ 
using  $\langle B \in \text{sets } M \rangle$   $\langle B \subseteq A \rangle$  by auto
thus ?thesis by auto
qed

lemma inf-neg-ge-1:
assumes A ∈ sets M
and  $\neg \text{pos-meas-set } A$ 
shows  $(1:\text{nat}) \leq \text{inf-neg } A$ 
proof –
define N where  $N = \{n:\text{nat} | n. (1:\text{nat}) \leq n \wedge$ 
 $(\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < \text{ereal}(-1/n))\}$ 
have  $N \neq \{\}$  unfolding N-def using assms inf-neg-ne by auto
moreover have  $\bigwedge n. n \in N \implies (1:\text{nat}) \leq n$  unfolding N-def by simp
ultimately show  $1 \leq \text{inf-neg } A$  unfolding inf-neg-def N-def
using Inf-nat-def1 assms(1) assms(2) by presburger
qed

lemma inf-neg-pos:
assumes A ∈ sets M
and  $\neg \text{pos-meas-set } A$ 
shows  $\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < -1/(\text{inf-neg } A)$ 
proof –
define N where  $N = \{n:\text{nat} | n. (1:\text{nat}) \leq n \wedge$ 
 $(\exists B \in \text{sets } M. B \subseteq A \wedge \mu B < \text{ereal}(-1/n))\}$ 
have  $N \neq \{\}$  unfolding N-def using assms inf-neg-ne by auto
hence Inf N ∈ N using Inf-nat-def1[of  $N$ ] by simp
hence  $\text{inf-neg } A \in N$  unfolding N-def inf-neg-def using assms by auto
thus ?thesis unfolding N-def by auto
qed

```

definition rep-neg **where**

rep-neg A = (if (A ∉ sets M ∨ pos-meas-set A) then {} else
SOME B. B ∈ sets M ∧ B ⊆ A ∧ μ B ≤ ereal (-1 / (inf-neg A)))

lemma g-rep-neg:

assumes A ∈ sets M

and ¬ pos-meas-set A

shows rep-neg A ∈ sets M rep-neg A ⊆ A

$\mu(\text{rep-neg } A) \leq \text{ereal}(-1 / (\text{inf-neg } A))$

proof –

have ∃ B. B ∈ sets M ∧ B ⊆ A ∧ μ B ≤ -1 / (inf-neg A) **using** assms

inf-neg-pos[of A] **by** auto

from someI-ex[OF this] **show** rep-neg A ∈ sets M rep-neg A ⊆ A

$\mu(\text{rep-neg } A) \leq -1 / (\text{inf-neg } A)$

unfolding rep-neg-def **using** assms **by** auto

qed

lemma rep-neg-sets:

shows rep-neg A ∈ sets M

proof (cases A ∉ sets M ∨ pos-meas-set A)

case True

then show ?thesis **unfolding** rep-neg-def **by** simp

next

case False

then show ?thesis **using** g-rep-neg(1) **by** blast

qed

lemma rep-neg-subset:

shows rep-neg A ⊆ A

proof (cases A ∉ sets M ∨ pos-meas-set A)

case True

then show ?thesis **unfolding** rep-neg-def **by** simp

next

case False

then show ?thesis **using** g-rep-neg(2) **by** blast

qed

lemma rep-neg-less:

assumes A ∈ sets M

and ¬ pos-meas-set A

shows $\mu(\text{rep-neg } A) \leq \text{ereal}(-1 / (\text{inf-neg } A))$ **using** assms g-rep-neg(3)

by simp

lemma rep-neg-leq:

shows $\mu(\text{rep-neg } A) \leq 0$

proof (cases A ∉ sets M ∨ pos-meas-set A)

case True

hence rep-neg A = {} **unfolding** rep-neg-def **by** simp

```

then show ?thesis using sgn-meas signed-measure-empty by force
next
  case False
  then show ?thesis using rep-neg-less by (metis le-ereal-le minus-divide-left
    neg-le-0-iff-le of-nat-0 of-nat-le-iff zero-ereal-def zero-le
    zero-le-divide-1-iff)
qed

```

3.2 Construction of the positive subset

```

fun pos-wtn
where
  pos-wtn-base: pos-wtn E 0 = E|
  pos-wtn-step: pos-wtn E (Suc n) = pos-wtn E n - rep-neg (pos-wtn E n)

lemma pos-wtn-subset:
  shows pos-wtn E n ⊆ E
  proof (induct n)
    case 0
    then show ?case using pos-wtn-base by simp
  next
    case (Suc n)
    hence rep-neg (pos-wtn E n) ⊆ pos-wtn E n using rep-neg-subset by simp
    then show ?case using Suc by auto
  qed

lemma pos-wtn-sets:
  assumes E ∈ sets M
  shows pos-wtn E n ∈ sets M
  proof (induct n)
    case 0
    then show ?case using assms by simp
  next
    case (Suc n)
    then show ?case using pos-wtn-step rep-neg-sets by auto
  qed

definition neg-wtn where
  neg-wtn E (n::nat) = rep-neg (pos-wtn E n)

lemma neg-wtn-neg-meas:
  shows μ (neg-wtn E n) ≤ 0 unfolding neg-wtn-def using rep-neg-leq by simp

lemma neg-wtn-sets:
  shows neg-wtn E n ∈ sets M unfolding neg-wtn-def using rep-neg-sets by simp

lemma neg-wtn-subset:
  shows neg-wtn E n ⊆ E unfolding neg-wtn-def
  using pos-wtn-subset[of E n] rep-neg-subset[of pos-wtn E n] by simp

```

```

lemma neg-wtn-union-subset:
  shows ( $\bigcup i \leq n. \text{neg-wtn } E i$ )  $\subseteq E$  using neg-wtn-subset by auto

lemma pos-wtn-Suc:
  shows pos-wtn E ( $\text{Suc } n$ ) =  $E - (\bigcup i \leq n. \text{neg-wtn } E i)$  unfolding neg-wtn-def
proof (induct n)
  case 0
  then show ?case using pos-wtn-base pos-wtn-step by simp
next
  case ( $\text{Suc } n$ )
  have pos-wtn E ( $\text{Suc } (\text{Suc } n)$ ) = pos-wtn E ( $\text{Suc } n$ ) -
    rep-neg (pos-wtn E ( $\text{Suc } n$ ))
    using pos-wtn-step by simp
  also have ... =  $(E - (\bigcup i \leq n. \text{rep-neg } (\text{pos-wtn } E i))) -$ 
    rep-neg (pos-wtn E ( $\text{Suc } n$ ))
    using Suc by simp
  also have ... =  $E - (\bigcup i \leq (\text{Suc } n). \text{rep-neg } (\text{pos-wtn } E i))$ 
    using diff-union[of E  $\lambda i. \text{rep-neg } (\text{pos-wtn } E i) n$ ] by auto
  finally show pos-wtn E ( $\text{Suc } (\text{Suc } n)$ ) =
     $E - (\bigcup i \leq (\text{Suc } n). \text{rep-neg } (\text{pos-wtn } E i))$  .
qed

definition pos-sub where
  pos-sub E = ( $\bigcap n. \text{pos-wtn } E n$ )

lemma pos-sub-sets:
  assumes E ∈ sets M
  shows pos-sub E ∈ sets M unfolding pos-sub-def using pos-wtn-sets assms
  by auto

lemma pos-sub-subset:
  shows pos-sub E  $\subseteq E$  unfolding pos-sub-def using pos-wtn-subset by blast

lemma pos-sub-infty:
  assumes E ∈ sets M
  and  $|\mu E| < \infty$ 
  shows  $|\mu (\text{pos-sub } E)| < \infty$  using signed-measure-finite-subset assms
  pos-sub-sets pos-sub-subset by simp

lemma neg-wtn-djn:
  shows disjoint-family ( $\lambda n. \text{neg-wtn } E n$ ) unfolding disjoint-family-on-def
proof -
  {
    fix n
    fix m::nat
    assume n < m
    hence  $\exists p. m = \text{Suc } p$  using old.nat.exhaust by auto
    from this obtain p where m = Suc p by auto
  }

```

```

have neg-wtn E m ⊆ pos-wtn E m unfolding neg-wtn-def
  by (simp add: rep-neg-subset)
also have ... = E - (⋃ i ≤ p. neg-wtn E i) using pos-wtn-Suc ⟨m = Suc p⟩
  by simp
finally have neg-wtn E m ⊆ E - (⋃ i ≤ p. neg-wtn E i) .
moreover have neg-wtn E n ⊆ (⋃ i ≤ p. neg-wtn E i) using ⟨n < m⟩
  ⟨m = Suc p⟩ by (simp add: UN-upper)
ultimately have neg-wtn E n ∩ neg-wtn E m = {} by auto
}
thus ∀ m ∈ UNIV. ∀ n ∈ UNIV. m ≠ n → neg-wtn E m ∩ neg-wtn E n = {}
  by (metis inf-commute linorder-neqE-nat)
qed
end

lemma disjoint-family-imp-on:
assumes disjoint-family A
shows disjoint-family-on A S
using assms disjoint-family-on-mono subset-UNIV by blast

context signed-measure-space
begin

lemma neg-wtn-union-neg-meas:
shows μ (⋃ i ≤ n. neg-wtn E i) ≤ 0
proof -
  have μ (⋃ i ≤ n. neg-wtn E i) = (∑ i ∈ {.. n}. μ (neg-wtn E i))
  proof (rule signed-measure-disj-sum, simp+)
    show signed-measure M μ using sgn-meas .
    show disjoint-family-on (neg-wtn E) {..n} using neg-wtn-djn
      disjoint-family-imp-on[neg-wtn E] by simp
    show ∀ i. i ∈ {..n} ⇒ neg-wtn E i ∈ sets M using neg-wtn-sets by simp
  qed
  also have ... ≤ 0 using neg-wtn-neg-meas by (simp add: sum-nonpos)
  finally show ?thesis .
qed

lemma pos-wtn-meas-gt:
assumes 0 < μ E
and E ∈ sets M
shows 0 < μ (pos-wtn E n)
proof (cases n = 0)
  case True
  then show ?thesis using assms by simp
next
  case False
  hence ∃ m. n = Suc m by (simp add: not0-implies-Suc)
  from this obtain m where n = Suc m by auto
  hence eq: pos-wtn E n = E - (⋃ i ≤ m. neg-wtn E i) using pos-wtn-Suc
    by simp
  hence pos-wtn E n ∩ (⋃ i ≤ m. neg-wtn E i) = {} by auto

```

```

moreover have  $E = \text{pos-wtn } E n \cup (\bigcup i \leq m. \text{neg-wtn } E i)$ 
  using eq neg-wtn-union-subset[of  $E m$ ] by auto
ultimately have  $\mu E = \mu(\text{pos-wtn } E n) + \mu(\bigcup i \leq m. \text{neg-wtn } E i)$ 
  using signed-measure-add[of  $M \mu \text{pos-wtn } E n \bigcup i \leq m. \text{neg-wtn } E i$ ]
    pos-wtn-sets neg-wtn-sets assms sgn-meas by auto
hence  $0 < \mu(\text{pos-wtn } E n) + \mu(\bigcup i \leq m. \text{neg-wtn } E i)$  using assms by simp
thus ?thesis using neg-wtn-union-neg-meas
  by (metis add.right-neutral add-mono not-le)
qed

definition union-wit where
  union-wit  $E = (\bigcup n. \text{neg-wtn } E n)$ 

lemma union-wit-sets:
  shows union-wit  $E \in \text{sets } M$  unfolding union-wit-def
proof (intro sigma-algebra.countable-nat-UN)
  show sigma-algebra (space  $M$ ) (sets  $M$ )
    by (simp add: sets.sigma-algebra-axioms)
  show range (neg-wtn  $E$ )  $\subseteq \text{sets } M$ 
  proof -
    {
      fix  $n$ 
      have neg-wtn  $E n \in \text{sets } M$  unfolding neg-wtn-def
        by (simp add: rep-neg-sets)
    }
    thus ?thesis by auto
  qed
qed

lemma union-wit-subset:
  shows union-wit  $E \subseteq E$ 
proof -
  {
    fix  $n$ 
    have neg-wtn  $E n \subseteq E$  unfolding neg-wtn-def using pos-wtn-subset
      rep-neg-subset[of pos-wtn  $E n$ ] by auto
  }
  thus ?thesis unfolding union-wit-def by auto
qed

lemma pos-sub-diff:
  shows pos-sub  $E = E - \text{union-wit } E$ 
proof
  show pos-sub  $E \subseteq E - \text{union-wit } E$ 
  proof -
    have pos-sub  $E \subseteq E$  using pos-sub-subset by simp
    moreover have pos-sub  $E \cap \text{union-wit } E = \{\}$ 
    proof (rule ccontr)
      assume pos-sub  $E \cap \text{union-wit } E \neq \{\}$ 

```

```

hence  $\exists a. a \in pos\text{-}sub E \cap union\text{-}wit E$  by auto
from this obtain a where  $a \in pos\text{-}sub E \cap union\text{-}wit E$  by auto
hence  $a \in union\text{-}wit E$  by simp
hence  $\exists n. a \in rep\text{-}neg (pos\text{-}wtn E n)$  unfolding union-wit-def neg-wtn-def
      by auto
from this obtain n where  $a \in rep\text{-}neg (pos\text{-}wtn E n)$  by auto
have  $a \in pos\text{-}wtn E (Suc n)$  using  $\langle a \in pos\text{-}sub E \cap union\text{-}wit E \rangle$ 
      unfolding pos-sub-def by blast
hence  $a \notin rep\text{-}neg (pos\text{-}wtn E n)$  using pos-wtn-step by simp
thus False using  $\langle a \in rep\text{-}neg (pos\text{-}wtn E n) \rangle$  by simp
qed
ultimately show ?thesis by auto
qed
next
show  $E - union\text{-}wit E \subseteq pos\text{-}sub E$ 
proof
fix a
assume  $a \in E - union\text{-}wit E$ 
show  $a \in pos\text{-}sub E$  unfolding pos-sub-def
proof
fix n
show  $a \in pos\text{-}wtn E n$ 
proof (cases  $n = 0$ )
case True
thus ?thesis using pos-wtn-base  $\langle a \in E - union\text{-}wit E \rangle$  by simp
next
case False
hence  $\exists m. n = Suc m$  by (simp add: not0-implies-Suc)
from this obtain m where  $n = Suc m$  by auto
have  $(\bigcup i \leq m. rep\text{-}neg (pos\text{-}wtn E i)) \subseteq$ 
       $(\bigcup n. (rep\text{-}neg (pos\text{-}wtn E n)))$  by auto
hence  $a \in E - (\bigcup i \leq m. rep\text{-}neg (pos\text{-}wtn E i))$ 
      using  $\langle a \in E - union\text{-}wit E \rangle$  unfolding union-wit-def neg-wtn-def
      by auto
thus  $a \in pos\text{-}wtn E n$  using pos-wtn-Suc  $\langle n = Suc m \rangle$ 
      unfolding neg-wtn-def by simp
qed
qed
qed
qed
qed

```

definition num-wtn **where**
 $num\text{-}wtn E n = inf\text{-}neg (pos\text{-}wtn E n)$

lemma num-wtn-geq:
shows $\mu (neg\text{-}wtn E n) \leq ereal (-1 / (num\text{-}wtn E n))$
proof (cases $(pos\text{-}wtn E n) \notin sets M \vee pos\text{-}meas\text{-}set (pos\text{-}wtn E n)$)
 case True
 hence $neg\text{-}wtn E n = \{\}$ unfolding neg-wtn-def rep-neg-def by simp

```

moreover have num-wtn E n = 0 using True unfolding num-wtn-def inf-neg-def
    by simp
ultimately show ?thesis using sgn-meas signed-measure-empty by force
next
case False
then show ?thesis using g-rep-neg(3)[of pos-wtn E n] unfolding neg-wtn-def
    num-wtn-def by simp
qed

lemma neg-wtn-infty:
assumes E ∈ sets M
and |μ E| < ∞
shows |μ (neg-wtn E i)| < ∞
proof (rule signed-measure-finite-subset)
show E ∈ sets M |μ E| < ∞ using assms by auto
show neg-wtn E i ∈ sets M
proof (cases pos-wtn E i ∉ sets M ∨ pos-meas-set (pos-wtn E i))
case True
then show ?thesis unfolding neg-wtn-def rep-neg-def by simp
next
case False
then show ?thesis unfolding neg-wtn-def
    using g-rep-neg(1)[of pos-wtn E i] by simp
qed
show neg-wtn E i ⊆ E unfolding neg-wtn-def using pos-wtn-subset[of E]
    rep-neg-subset[of pos-wtn E i] by auto
qed

lemma union-wit-infty:
assumes E ∈ sets M
and |μ E| < ∞
shows |μ (union-wit E)| < ∞ using union-wit-subset union-wit-sets
    signed-measure-finite-subset assms unfolding union-wit-def by simp

lemma neg-wtn-summable:
assumes E ∈ sets M
and |μ E| < ∞
shows summable (λi. - real-of-ereal (μ (neg-wtn E i)))
proof -
have signed-measure M μ using sgn-meas .
moreover have range (neg-wtn E) ⊆ sets M unfolding neg-wtn-def
    using rep-neg-sets by auto
moreover have disjoint-family (neg-wtn E) using neg-wtn-djn by simp
moreover have ∪ (range (neg-wtn E)) ∈ sets M using union-wit-sets
    unfolding union-wit-def by simp
moreover have |μ (∪ (range (neg-wtn E)))| < ∞
    using union-wit-subset signed-measure-finite-subset union-wit-sets assms
    unfolding union-wit-def by simp

```

ultimately have summable ($\lambda i. \text{real-of-ereal } |\mu(\text{neg-wtn } E i)|$)
 using signed-measure-abs-convergent[of M] by simp
 moreover have $\bigwedge i. |\mu(\text{neg-wtn } E i)| = -(\mu(\text{neg-wtn } E i))$
proof –
 fix i
 have $\mu(\text{neg-wtn } E i) \leq 0$ using rep-neg-leq[of pos-wtn $E i$]
 unfolding neg-wtn-def .
 thus $|\mu(\text{neg-wtn } E i)| = -\mu(\text{neg-wtn } E i)$ using less-eq-ereal-def by auto
qed
 ultimately show ?thesis by simp
qed

lemma inv-num-wtn-summable:

assumes $E \in \text{sets } M$
 and $|\mu E| < \infty$
shows summable ($\lambda n. 1/(\text{num-wtn } E n)$)
proof (rule summable-bounded)
 show $\bigwedge i. 0 \leq 1 / \text{real } (\text{num-wtn } E i)$ by simp
 show $\bigwedge i. 1 / \text{real } (\text{num-wtn } E i) \leq (\lambda n. -\text{real-of-ereal } (\mu(\text{neg-wtn } E n))) i$
proof –
 fix i
 have $|\mu(\text{neg-wtn } E i)| < \infty$ using assms neg-wtn-infty by simp
 have $\text{ereal } (1/(\text{num-wtn } E i)) \leq -\mu(\text{neg-wtn } E i)$ using num-wtn-geq[of $E i$]
 ereal-minus-le-minus by fastforce
 also have ... = $\text{ereal}(-\text{real-of-ereal } (\mu(\text{neg-wtn } E i)))$
 using $\langle |\mu(\text{neg-wtn } E i)| < \infty \rangle$ ereal-real' by auto
 finally have $\text{ereal } (1/(\text{num-wtn } E i)) \leq$
 $\text{ereal}(-\text{real-of-ereal } (\mu(\text{neg-wtn } E i)))$.
 thus $1 / \text{real } (\text{num-wtn } E i) \leq -\text{real-of-ereal } (\mu(\text{neg-wtn } E i))$ by simp
qed
 show summable ($\lambda i. -\text{real-of-ereal } (\mu(\text{neg-wtn } E i))$)
 using assms neg-wtn-summable by simp
qed

lemma inv-num-wtn-shift-summable:

assumes $E \in \text{sets } M$
 and $|\mu E| < \infty$
shows summable ($\lambda n. 1/(\text{num-wtn } E n - 1)$)
proof (rule sum-shift-denum)
 show summable ($\lambda n. 1 / \text{real } (\text{num-wtn } E n)$) using assms inv-num-wtn-summable
 by simp
qed

lemma neg-wtn-meas-sums:

assumes $E \in \text{sets } M$
 and $|\mu E| < \infty$
shows ($\lambda i. -(\mu(\text{neg-wtn } E i))$) sums
 $\text{suminf } (\lambda i. -\text{real-of-ereal } (\mu(\text{neg-wtn } E i)))$
proof –

```

have (λi. ereal (− real-of-ereal (μ (neg-wtn E i)))) sums
  suminf (λi. − real-of-ereal (μ (neg-wtn E i)))
proof (rule sums-ereal[THEN iffD2])
  have summable (λi. − real-of-ereal (μ (neg-wtn E i)))
    using neg-wtn-summable assms by simp
  thus (λx. − real-of-ereal (μ (neg-wtn E x)))
    sums (Σ i. − real-of-ereal (μ (neg-wtn E i)))
    by auto
  qed
  moreover have ∀i. μ (neg-wtn E i) = ereal (real-of-ereal (μ (neg-wtn E i)))
  proof –
    fix i
    show μ (neg-wtn E i) = ereal (real-of-ereal (μ (neg-wtn E i)))
      using assms(1) assms(2) ereal-real' neg-wtn-inf by auto
    qed
    ultimately show ?thesis
      by (metis (no-types, lifting) sums-cong uminus-ereal.simps(1))
  qed

lemma neg-wtn-meas-suminf-le:
  assumes E ∈ sets M
  and |μ E| < ∞
  shows suminf (λi. μ (neg-wtn E i)) ≤ − suminf (λn. 1/(num-wtn E n))
  proof –
    have suminf (λn. 1/(num-wtn E n)) ≤
      suminf (λi. −real-of-ereal (μ (neg-wtn E i)))
    proof (rule suminf-le)
      show summable (λn. 1 / real (num-wtn E n)) using assms
        inv-num-wtn-summable[of E]
        summable-minus[of λn. 1 / real (num-wtn E n)] by simp
      show summable (λi. −real-of-ereal (μ (neg-wtn E i)))
        using neg-wtn-summable assms
          summable-minus[of λi. real-of-ereal (μ (neg-wtn E i))]
        by (simp add: summable-minus-iff)
      show ∀n. 1 / real (num-wtn E n) ≤ −real-of-ereal (μ (neg-wtn E n))
      proof –
        fix n
        have μ (neg-wtn E n) ≤ ereal (− 1 / real (num-wtn E n))
          using num-wtn-geq by simp
        hence ereal (1 / real (num-wtn E n)) ≤ − μ (neg-wtn E n)
          by (metis add.inverse-inverse eq-iff ereal-uminus-le-reorder linear
              minus-divide-left uminus-ereal.simps(1))
        have real-of-ereal (ereal (1 / real (num-wtn E n))) ≤
          real-of-ereal (− μ (neg-wtn E n))
        proof (rule real-of-ereal-positive-mono)
          show 0 ≤ ereal (1 / real (num-wtn E n)) by simp
          show ereal (1 / real (num-wtn E n)) ≤ − μ (neg-wtn E n)
            using ereal (1 / real (num-wtn E n)) ≤ − μ (neg-wtn E n) .
          show − μ (neg-wtn E n) ≠ ∞ using neg-wtn-inf by auto
    qed
  qed

```

```

qed
thus  $(1 / \text{real}(\text{num-wtn } E n)) \leq -\text{real-of-ereal}(\mu(\text{neg-wtn } E n))$ 
  by simp
qed
qed
also have ... =  $-\text{suminf}(\lambda i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i)))$ 
proof (rule suminf-minus)
  show summable  $(\lambda n. \text{real-of-ereal}(\mu(\text{neg-wtn } E n)))$ 
    using neg-wtn-summable assms
      sumable-minus[of  $\lambda i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i))]$ 
    by (simp add: summable-minus-iff)
qed
finally have suminf  $(\lambda n. 1/(\text{num-wtn } E n)) \leq$ 
   $-\text{suminf}(\lambda i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i)))$ .
hence a: suminf  $(\lambda i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i))) \leq$ 
   $-\text{suminf}(\lambda n. 1/(\text{num-wtn } E n))$  by simp
show suminf  $(\lambda i. (\mu(\text{neg-wtn } E i))) \leq \text{ereal}(-\text{suminf}(\lambda n. 1/(\text{num-wtn } E n)))$ 

proof -
  have sumeq: suminf  $(\lambda i. \text{ereal}(\text{real-of-ereal}(\mu(\text{neg-wtn } E i)))) =$ 
    suminf  $(\lambda i. (\text{real-of-ereal}(\mu(\text{neg-wtn } E i))))$ 
proof (rule sums-suminf-ereal)
  have summable  $(\lambda i. -\text{real-of-ereal}(\mu(\text{neg-wtn } E i)))$ 
    using neg-wtn-summable assms
      sumable-minus[of  $\lambda i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i))]$ 
    by (simp add: summable-minus-iff)
  thus  $(\lambda i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i))) \text{ sums}$ 
     $(\sum i. \text{real-of-ereal}(\mu(\text{neg-wtn } E i)))$ 
    using neg-wtn-summable[of E] assms sumable-minus-iff by blast
qed
hence suminf  $(\lambda i. \mu(\text{neg-wtn } E i)) =$ 
  suminf  $(\lambda i. (\text{real-of-ereal}(\mu(\text{neg-wtn } E i))))$ 
proof -
  have  $\bigwedge i. \text{ereal}(\text{real-of-ereal}(\mu(\text{neg-wtn } E i))) = \mu(\text{neg-wtn } E i)$ 
  proof -
    fix i
    show ereal  $(\text{real-of-ereal}(\mu(\text{neg-wtn } E i))) = \mu(\text{neg-wtn } E i)$ 
      using neg-wtn-infty[of E] assms by (simp add: ereal-real')
  qed
  thus ?thesis using sumeq by auto
qed
thus ?thesis using a by simp
qed
qed

```

lemma union-wit-meas-le:
assumes $E \in \text{sets } M$
and $|\mu E| < \infty$
shows $\mu(\text{union-wit } E) \leq -\text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n))$

```

proof -
  have  $\mu(\text{union-wit } E) = \mu(\bigcup (\text{range}(\text{neg-wtn } E)))$  unfolding union-wit-def
    by simp
  also have ... =  $(\sum i. \mu(\text{neg-wtn } E i))$ 
  proof (rule signed-measure-inf-sum[symmetric])
    show signed-measure  $M$   $\mu$  using sgn-meas .
    show range( $\text{neg-wtn } E$ )  $\subseteq$  sets  $M$ 
      by (simp add: image-subset-iff neg-wtn-def rep-neg-sets)
    show disjoint-family( $\text{neg-wtn } E$ ) using neg-wtn-djn by simp
    show  $\bigcup (\text{range}(\text{neg-wtn } E)) \in \text{sets } M$  using union-wit-sets
      unfolding union-wit-def by simp
  qed
  also have ...  $\leq -\text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n))$ 
    using assms neg-wtn-meas-suminf-le by simp
    finally show ?thesis .
  qed

lemma pos-sub-pos-meas:
  assumes  $E \in \text{sets } M$ 
  and  $|\mu E| < \infty$ 
  and  $0 < \mu E$ 
  and  $\neg \text{pos-meas-set } E$ 
  shows  $0 < \mu(\text{pos-sub } E)$ 
  proof -
    have  $0 < \mu E$  using assms by simp
    also have ... =  $\mu(\text{pos-sub } E) + \mu(\text{union-wit } E)$ 
    proof -
      have  $E = \text{pos-sub } E \cup (\text{union-wit } E)$ 
        using pos-sub-diff[of  $E$ ] union-wit-subset by force
      moreover have  $\text{pos-sub } E \cap \text{union-wit } E = \{\}$ 
        using pos-sub-diff by auto
      ultimately show ?thesis
        using signed-measure-add[of  $M$   $\mu$  pos-sub  $E$  union-wit  $E$ ]
          pos-sub-sets union-wit-sets assms sgn-meas by simp
    qed
    also have ...  $\leq \mu(\text{pos-sub } E) + (-\text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n)))$ 
    proof -
      have  $\mu(\text{union-wit } E) \leq -\text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n))$ 
        using union-wit-meas-le[of  $E$ ] assms by simp
      thus ?thesis using union-wit-infty assms using add-left-mono by blast
    qed
    also have ... =  $\mu(\text{pos-sub } E) - \text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n))$ 
      by (simp add: minus-ereal-def)
    finally have  $0 < \mu(\text{pos-sub } E) - \text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n))$  .
    moreover have  $0 < \text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n))$ 
    proof (rule suminf-pos2)
      show  $0 < 1 / \text{real}(\text{num-wtn } E 0)$ 
        using inf-neg-ge-1[of  $E$ ] assms pos-wtn-base unfolding num-wtn-def by simp
      show  $\bigwedge n. 0 \leq 1 / \text{real}(\text{num-wtn } E n)$  by simp
  
```

```

show summable ( $\lambda n. 1 / \text{real}(\text{num-wtn } E n)$ )
  using assms inv-num-wtn-summable by simp
qed
ultimately show ?thesis using pos-sub-inf $\ell$ ty assms by fastforce

lemma num-wtn-conv:
assumes E ∈ sets M
  and |μ E| < ∞
shows ( $\lambda n. 1 / (\text{num-wtn } E n)$ ) —→ 0
proof (rule summable-LIMSEQ-zero)
  show summable ( $\lambda n. 1 / \text{real}(\text{num-wtn } E n)$ )
    using assms inv-num-wtn-summable by simp
qed

lemma num-wtn-shift-conv:
assumes E ∈ sets M
  and |μ E| < ∞
shows ( $\lambda n. 1 / (\text{num-wtn } E n - 1)$ ) —→ 0
proof (rule summable-LIMSEQ-zero)
  show summable ( $\lambda n. 1 / \text{real}(\text{num-wtn } E n - 1)$ )
    using assms inv-num-wtn-shift-summable by simp
qed

lemma inf-neg-E-set:
assumes 0 < inf-neg E
shows E ∈ sets M using assms unfolding inf-neg-def by presburger

lemma inf-neg-pos-meas:
assumes 0 < inf-neg E
shows ¬ pos-meas-set E using assms unfolding inf-neg-def by presburger

lemma inf-neg-mem:
assumes 0 < inf-neg E
shows inf-neg E ∈ {n::nat | n. (1::nat) ≤ n ∧
  (exists B ∈ sets M. B ⊆ E ∧ μ B < ereal (-1/n))}

proof -
  have E ∈ sets M using assms unfolding inf-neg-def by presburger
  moreover have ¬ pos-meas-set E using assms unfolding inf-neg-def
    by presburger
  ultimately have {n::nat | n. (1::nat) ≤ n ∧
    (exists B ∈ sets M. B ⊆ E ∧ μ B < ereal (-1/n))} ≠ {}
    using inf-neg-ne[of E] by simp
  thus ?thesis unfolding inf-neg-def
    by (meson Inf-nat-def1 ‹E ∈ sets M› ‹¬ pos-meas-set E›)
qed

lemma prec-inf-neg-pos:
assumes 0 < inf-neg E - 1

```

and $B \in \text{sets } M$
 and $B \subseteq E$
 shows $-1/(\inf\text{-}neg E - 1) \leq \mu B$
proof (rule *ccontr*)
 define S where $S = \{p::nat | p. (1::nat) \leq p \wedge (\exists B \in \text{sets } M. B \subseteq E \wedge \mu B < \text{ereal } (-1/p))\}$
 assume $\neg \text{ereal } (-1 / \text{real } (\inf\text{-}neg E - 1)) \leq \mu B$
 hence $\mu B < -1/(\inf\text{-}neg E - 1)$ by *auto*
 hence $\inf\text{-}neg E - 1 \in S$ unfolding $S\text{-def}$ using *assms* by *auto*
 have $\text{Suc } 0 < \inf\text{-}neg E$ using *assms* by *simp*
 hence $\inf\text{-}neg E \in S$ unfolding $S\text{-def}$ using *inf-neg-mem[of E]* by *simp*
 hence $S \neq \{\}$ by *auto*
 have $\inf\text{-}neg E = \text{Inf } S$ unfolding $S\text{-def}$ *inf-neg-def*
 using *assms* *inf-neg-E-set inf-neg-pos-meas* by *auto*
 have $\inf\text{-}neg E - 1 < \inf\text{-}neg E$ using *assms* by *simp*
 hence $\inf\text{-}neg E - 1 \notin S$
 using *cInf-less-iff[of S] {S} \neq \{\} \langle \inf\text{-}neg E = \text{Inf } S \rangle* by *auto*
 thus False using $\langle \inf\text{-}neg E - 1 \in S \rangle$ by *simp*
qed

lemma *pos-wtn-meas-ge*:
 assumes $E \in \text{sets } M$
 and $|\mu E| < \infty$
 and $C \in \text{sets } M$
 and $\bigwedge_n. C \subseteq \text{pos-wtn } E n$
 and $\bigwedge_n. 0 < \text{num-wtn } E n$
 shows $\exists N. \forall n \geq N. -1/(\text{num-wtn } E n - 1) \leq \mu C$
proof –
 have $\exists N. \forall n \geq N. 1/(\text{num-wtn } E n) < 1/2$ using *num-wtn-conv[of E]*
 conv-0-half[of λn. 1 / real (num-wtn E n)] assms by *simp*
 from this obtain N where $\forall n \geq N. 1/(\text{num-wtn } E n) < 1/2$ by *auto*
 {
 fix n
 assume $N \leq n$
 hence $1/(\text{num-wtn } E n) < 1/2$ using $\langle \forall n \geq N. 1/(\text{num-wtn } E n) < 1/2 \rangle$ by
 simp
 have $1/(1/2) < 1/(1/(\text{num-wtn } E n))$
 proof (rule *frac-less2*, *auto*)
 show $2 / \text{real } (\text{num-wtn } E n) < 1$ using $\langle 1/(\text{num-wtn } E n) < 1/2 \rangle$
 by *linarith*
 show $0 < \text{num-wtn } E n$ unfolding *num-wtn-def* using *inf-neg-ge-1 assms*
 by (*simp add: num-wtn-def*)
 qed
 hence $2 < (\text{num-wtn } E n)$ by *simp*
 hence $\text{Suc } 0 < \text{num-wtn } E n - 1$ unfolding *num-wtn-def* by *simp*
 hence $-1/(\text{num-wtn } E n - 1) \leq \mu C$ using *assms prec-inf-neg-pos*
 unfolding *num-wtn-def* by *simp*
 }
 thus *?thesis* by *auto*

qed

lemma pos-sub-pos-meas-subset:

assumes $E \in \text{sets } M$

and $|\mu E| < \infty$

and $C \in \text{sets } M$

and $C \subseteq (\text{pos-sub } E)$

and $\bigwedge n. 0 < \text{num-wtn } E n$

shows $0 \leq \mu C$

proof –

have $\bigwedge n. C \subseteq \text{pos-wtn } E n$ using assms unfolding pos-sub-def by auto

hence $\exists N. \forall n \geq N. -1 / (\text{num-wtn } E n - 1) \leq \mu C$ using assms

pos-wtn-meas-ge[of $E C$] by simp

from this obtain N where $Nprop: \forall n \geq N. -1 / (\text{num-wtn } E n - 1) \leq \mu C$

by auto

show $0 \leq \mu C$

proof (rule lim-mono)

show $\bigwedge n. N \leq n \implies -1 / (\text{num-wtn } E n - 1) \leq (\lambda n. \mu C) n$

using $Nprop$ by simp

have $(\lambda n. (-1 / \text{real}(\text{num-wtn } E n - 1))) \longrightarrow 0$

using assms num-wtn-shift-conv[of E] by simp

hence $(\lambda n. (-1 / \text{real}(\text{num-wtn } E n - 1))) \longrightarrow 0$

using tendsto-minus[of $\lambda n. 1 / \text{real}(\text{num-wtn } E n - 1) 0$] by simp

thus $(\lambda n. ereal(-1 / \text{real}(\text{num-wtn } E n - 1))) \longrightarrow 0$

by (simp add: zero-ereal-def)

show $(\lambda n. \mu C) \longrightarrow \mu C$ by simp

qed

qed

lemma pos-sub-pos-meas':

assumes $E \in \text{sets } M$

and $|\mu E| < \infty$

and $0 < \mu E$

and $\forall n. 0 < \text{num-wtn } E n$

shows $0 < \mu (\text{pos-sub } E)$

proof –

have $0 < \mu E$ using assms by simp

also have ... = $\mu (\text{pos-sub } E) + \mu (\text{union-wit } E)$

proof –

have $E = \text{pos-sub } E \cup (\text{union-wit } E)$

using pos-sub-diff[of E] union-wit-subset by force

moreover have $\text{pos-sub } E \cap \text{union-wit } E = \{\}$

using pos-sub-diff by auto

ultimately show ?thesis

using signed-measure-add[of $M \mu \text{pos-sub } E \text{union-wit } E$]

pos-sub-sets union-wit-sets assms sgn-meas by simp

qed

also have ... $\leq \mu (\text{pos-sub } E) + (-\text{suminf}(\lambda n. 1 / \text{real}(\text{num-wtn } E n)))$

proof –

```

have  $\mu(\text{union-wit } E) \leq -\text{suminf } (\lambda n. 1 / \text{real } (\text{num-wtn } E n))$ 
  using  $\text{union-wit-meas-le}[\text{of } E]$  assms by simp
thus ?thesis using  $\text{union-wit-infty}$  assms using  $\text{add-left-mono}$  by blast
qed
also have ... =  $\mu(\text{pos-sub } E) - \text{suminf } (\lambda n. 1 / \text{real } (\text{num-wtn } E n))$ 
  by (simp add: minus-ereal-def)
finally have  $0 < \mu(\text{pos-sub } E) - \text{suminf } (\lambda n. 1 / \text{real } (\text{num-wtn } E n))$  .
moreover have  $0 < \text{suminf } (\lambda n. 1 / \text{real } (\text{num-wtn } E n))$ 
proof (rule suminf-pos2)
  show  $0 < 1 / \text{real } (\text{num-wtn } E 0)$  using assms by simp
  show  $\bigwedge n. 0 \leq 1 / \text{real } (\text{num-wtn } E n)$  by simp
  show  $\text{summable } (\lambda n. 1 / \text{real } (\text{num-wtn } E n))$ 
    using assms inv-num-wtn-summable by simp
qed
ultimately show ?thesis using pos-sub-infty assms by fastforce
qed

```

We obtain the main result of this part on the existence of a positive subset.

```

lemma exists-pos-meas-subset:
assumes E ∈ sets M
and |μ E| < ∞
and 0 < μ E
shows ∃ A. A ⊆ E ∧ pos-meas-set A ∧ 0 < μ A
proof (cases ∀ n. 0 < num-wtn E n)
  case True
  have pos-meas-set (pos-sub E)
  proof (rule pos-meas-setI)
    show pos-sub E ∈ sets M by (simp add: assms(1) pos-sub-sets)
    fix A
    assume A ∈ sets M and A ⊆ pos-sub E
    thus 0 ≤ μ A using assms True pos-sub-pos-meas-subset[of E] by simp
  qed
  moreover have 0 < μ (pos-sub E)
    using pos-sub-pos-meas'[of E] True assms by simp
  ultimately show ?thesis using pos-meas-set-def by (metis pos-sub-subset)
next
  case False
  hence ∃ n. num-wtn E n = 0 by simp
  from this obtain n where num-wtn E n = 0 by auto
  hence pos-wtn E n ∉ sets M ∨ pos-meas-set (pos-wtn E n)
    using inf-neg-ge-1 unfolding num-wtn-def by fastforce
  hence pos-meas-set (pos-wtn E n) using assms
    by (simp add: ‹E ∈ sets M› pos-wtn-sets)
  moreover have 0 < μ (pos-wtn E n) using pos-wtn-meas-gt assms by simp
  ultimately show ?thesis using pos-meas-set-def by (meson pos-wtn-subset)
qed

```

4 The Hahn decomposition theorem

```

definition seq-meas where
  seq-meas = (SOME f. incseq f ∧ range f ⊆ pos-img ∧ ∐ pos-img = ∐ range f)

lemma seq-meas-props:
  shows incseq seq-meas ∧ range seq-meas ⊆ pos-img ∧
    ∐ pos-img = ∐ range seq-meas
proof –
  have ex: ∃f. incseq f ∧ range f ⊆ pos-img ∧ ∐ pos-img = ∐ range f
  proof (rule Extended-Real.Sup-countable-SUP)
  show pos-img ≠ {}
  proof –
    have {} ∈ pos-sets using empty-pos-meas-set unfolding pos-sets-def
    by simp
    hence μ {} ∈ pos-img unfolding pos-img-def by auto
    thus ?thesis by auto
  qed
qed
let ?V = SOME f. incseq f ∧ range f ⊆ pos-img ∧ ∐ pos-img = ∐ range f
have vprop: incseq ?V ∧ range ?V ⊆ pos-img ∧ ∐ pos-img = ∐ range ?V
  using someI-ex[of λf. incseq f ∧ range f ⊆ pos-img ∧
    ∐ pos-img = ∐ range f] ex by blast
show ?thesis using seq-meas-def vprop by presburger
qed

definition seq-meas-rep where
  seq-meas-rep n = (SOME A. A ∈ pos-sets ∧ seq-meas n = μ A)

lemma seq-meas-rep-ex:
  shows seq-meas-rep n ∈ pos-sets ∧ μ (seq-meas-rep n) = seq-meas n
proof –
  have ex: ∃A. A ∈ pos-sets ∧ seq-meas n = μ A using seq-meas-props
    by (smt (verit) UNIV-I image-subset-iff mem-Collect-eq pos-img-def)
  let ?V = SOME A. A ∈ pos-sets ∧ seq-meas n = μ A
  have vprop: ?V ∈ pos-sets ∧ seq-meas n = μ ?V using
    someI-ex[of λA. A ∈ pos-sets ∧ seq-meas n = μ A] using ex by blast
  show ?thesis using seq-meas-rep-def vprop by fastforce
qed

lemma seq-meas-rep-pos:
  assumes ∀E ∈ sets M. μ E < ∞
  shows pos-meas-set (∐ i. seq-meas-rep i)
proof (rule pos-meas-set-Union)
  show ∨i. pos-meas-set (seq-meas-rep i)
  using seq-meas-rep-ex signed-measure-space.pos-sets-def
    signed-measure-space-axioms by auto
  then show ∨i. seq-meas-rep i ∈ sets M
  by (simp add: pos-meas-setD1)

```

```

show |μ (⋃ (range seq-meas-rep))| < ∞
proof -
  have (⋃ (range seq-meas-rep)) ∈ sets M
  proof (rule sigma-algebra.countable-Union)
    show sigma-algebra (space M) (sets M)
      by (simp add: sets.sigma-algebra-axioms)
    show countable (range seq-meas-rep) by simp
    show range seq-meas-rep ⊆ sets M
      by (simp add: ⋀ i. seq-meas-rep i ∈ sets M image-subset-iff)
  qed
  hence μ (⋃ (range seq-meas-rep)) ≥ 0
    using ⋀ i. pos-meas-set (seq-meas-rep i) ⋀ i. seq-meas-rep i ∈ sets M
      signed-measure-space.pos-meas-set-pos-lim signed-measure-space-axioms
    by blast
  thus ?thesis using assms ⋃ (range seq-meas-rep) ∈ sets M abs-ereal-ge0
    by simp
  qed
qed

lemma sup-seq-meas-rep:
assumes ∀ E ∈ sets M. μ E < ∞
and S = (⋃ pos-img)
and A = (⋃ i. seq-meas-rep i)
shows μ A = S
proof -
  have pms: pos-meas-set (⋃ i. seq-meas-rep i)
    using assms seq-meas-rep-pos by simp
  hence μ A ≤ S
    by (metis (mono-tags, lifting) Sup-upper ⋃ S = ⋃ pos-img mem-Collect-eq
      pos-img-def pos-meas-setD1 pos-sets-def assms(2) assms(3))
  have ∀ n. (μ A = μ (A - seq-meas-rep n) + μ (seq-meas-rep n))
  proof
    fix n
    have A = (A - seq-meas-rep n) ∪ seq-meas-rep n
      using ⋃ A = ⋃ (range seq-meas-rep) by blast
    hence μ A = μ ((A - seq-meas-rep n) ∪ seq-meas-rep n) by simp
    also have ... = μ (A - seq-meas-rep n) + μ (seq-meas-rep n)
    proof (rule signed-measure-add)
      show signed-measure M μ using sgn-meas by simp
      show seq-meas-rep n ∈ sets M
        using pos-sets-def seq-meas-rep-ex by auto
      then show A - seq-meas-rep n ∈ sets M
        by (simp add: assms pms pos-meas-setD1 sets.Diff)
      show (A - seq-meas-rep n) ∩ seq-meas-rep n = {} by auto
    qed
    finally show μ A = μ (A - seq-meas-rep n) + μ (seq-meas-rep n).
  qed
  have ∀ n. μ A ≥ μ (seq-meas-rep n)
  proof

```

```

fix n
have  $\mu A \geq 0$  using pms assms unfolding pos-meas-set-def by auto
have  $(A - \text{seq-meas-rep } n) \subseteq A$  by simp
hence pos-meas-set  $(A - \text{seq-meas-rep } n)$ 
proof -
  have  $(A - \text{seq-meas-rep } n) \in \text{sets } M$ 
    using pms assms pos-meas-setD1 pos-sets-def seq-meas-rep-ex by auto
    thus ?thesis using pms assms unfolding pos-meas-set-def by auto
qed
hence  $\mu (A - \text{seq-meas-rep } n) \geq 0$  unfolding pos-meas-set-def by auto
thus  $\mu (\text{seq-meas-rep } n) \leq \mu A$ 
  using  $\forall n. (\mu A = \mu (A - \text{seq-meas-rep } n) + \mu (\text{seq-meas-rep } n))$ 
  by (metis erreal-le-add-self2)
qed
hence  $\mu A \geq (\bigsqcup \text{range seq-meas})$  by (simp add: Sup-le-iff seq-meas-rep-ex)
moreover have  $S = (\bigsqcup \text{range seq-meas})$ 
  using seq-meas-props  $\langle S = (\bigsqcup \text{pos-img}) \rangle$  by simp
ultimately have  $\mu A \geq S$  by simp
thus  $\mu A = S$  using  $\langle \mu A \leq S \rangle$  by simp
qed

lemma seq-meas-rep-compl:
assumes  $\forall E \in \text{sets } M. \mu E < \infty$ 
and  $A = (\bigcup i. \text{seq-meas-rep } i)$ 
shows neg-meas-set  $((\text{space } M) - A)$  unfolding neg-meas-set-def
proof (rule ccontr)
assume asm:  $\neg (\text{space } M - A \in \text{sets } M \wedge$ 
 $(\forall Aa \in \text{sets } M. Aa \subseteq \text{space } M - A \longrightarrow \mu Aa \leq 0))$ 
define S where  $S = (\bigsqcup \text{pos-img})$ 
have pos-meas-set A using assms seq-meas-rep-pos by simp
have  $\mu A = S$  using sup-seq-meas-rep assms S-def by simp
hence  $S < \infty$  using assms pos-meas-set A pos-meas-setD1 by blast
have  $(\text{space } M - A \in \text{sets } M)$ 
  by (simp add: pos-meas-set A pos-meas-setD1 sets.compl-sets)
hence  $\neg(\forall Aa \in \text{sets } M. Aa \subseteq \text{space } M - A \longrightarrow \mu Aa \leq 0)$  using asm by blast
hence  $\exists E \in \text{sets } M. E \subseteq ((\text{space } M) - A) \wedge \mu E > 0$ 
  by (metis less-eq-ereal-def linear)
from this obtain E where E in sets M and E ⊆ ((space M) - A) and
 $\mu E > 0$  by auto
have  $\exists A0 \subseteq E. \text{pos-meas-set } A0 \wedge \mu A0 > 0$ 
proof (rule exists-pos-meas-subset)
  show  $E \in \text{sets } M$  using  $\langle E \in \text{sets } M \rangle$  by simp
  show  $0 < \mu E$  using  $\langle \mu E > 0 \rangle$  by simp
  show  $|\mu E| < \infty$ 
proof -
  have  $\mu E < \infty$  using assms  $\langle E \in \text{sets } M \rangle$  by simp
  moreover have  $-\infty < \mu E$  using  $\langle 0 < \mu E \rangle$  by simp
  ultimately show ?thesis
  by (meson erreal-infty-less(1) not-inftyI)

```

```

qed
qed
from this obtain A0 where A0 ⊆ E and pos-meas-set A0 and μ A0 > 0
by auto
have pos-meas-set (A ∪ A0)
using pos-meas-set-union ⟨pos-meas-set A0⟩ ⟨pos-meas-set A⟩ by simp
have μ (A ∪ A0) = μ A + μ A0
proof (rule signed-measure-add)
show signed-measure M μ using sgn-meas by simp
show A ∈ sets M using ⟨pos-meas-set A⟩
unfolding pos-meas-set-def by simp
show A0 ∈ sets M using ⟨pos-meas-set A0⟩
unfolding pos-meas-set-def by simp
show (A ∩ A0) = {} using ⟨A0 ⊆ E⟩ ⟨E ⊆ ((space M) − A)⟩ by auto
qed
then have μ (A ∪ A0) > S
using ⟨μ A = S⟩ ⟨μ A0 > 0⟩
by (metis ‹S < ∞› ⟨pos-meas-set (A ∪ A0)⟩ abs-ereal-ge0 ereal-between(2)
not-inftyI not-less-iff-gr-or-eq pos-meas-self)
have (A ∪ A0) ∈ pos-sets
proof –
have (A ∪ A0) ∈ sets M using sigma-algebra.countable-Union
by (simp add: ⟨pos-meas-set (A ∪ A0)⟩ pos-meas-setD1)
moreover have pos-meas-set (A ∪ A0) using ⟨pos-meas-set (A ∪ A0)⟩ by
simp
ultimately show ?thesis unfolding pos-sets-def by simp
qed
then have μ (A ∪ A0) ∈ pos-img unfolding pos-img-def by auto
show False using ⟨μ (A ∪ A0) > S⟩ ⟨μ (A ∪ A0) ∈ pos-img⟩ ⟨S = (⊔ pos-img)⟩
by (metis Sup-upper sup.absorb-iff2 sup.strict-order-iff)
qed

lemma hahn-decomp-finite:
assumes ∀ E ∈ sets M. μ E < ∞
shows ∃ M1 M2. hahn-space-decomp M1 M2 unfolding hahn-space-decomp-def
proof –
define S where S = (⊔ pos-img)
define A where A = (⊔ i. seq-meas-rep i)
have pos-meas-set A unfolding A-def using assms seq-meas-rep-pos by simp
have neg-meas-set ((space M) − A)
using seq-meas-rep-compl assms unfolding A-def by simp
show ∃ M1 M2. pos-meas-set M1 ∧ neg-meas-set M2 ∧ space M = M1 ∪ M2 ∧
M1 ∩ M2 = {}
proof (intro exI conjI)
show pos-meas-set A using ⟨pos-meas-set A⟩ .
show neg-meas-set (space M − A) using ⟨neg-meas-set (space M − A)⟩ .
show space M = A ∪ (space M − A)
by (metis Diff-partition ⟨pos-meas-set A⟩ inf.absorb-iff2 pos-meas-setD1)

```

```

sets.Int-space-eq1)
show A ∩ (space M - A) = {} by auto
qed
qed

theorem hahn-decomposition:
shows ∃ M1 M2. hahn-space-decomp M1 M2
proof (cases ∀ E ∈ sets M. μ E < ∞)
case True
thus ?thesis using hahn-decomp-finite by simp
next
case False
define m where m = (λA . − μ A)
have ∃ M1 M2. signed-measure-space.hahn-space-decomp M m M1 M2
proof (rule signed-measure-space.hahn-decomp-finite)
show signed-measure-space M m
using signed-measure-minus sgn-meas ⟨m = (λA . − μ A)⟩
by (unfold-locales, simp)
show ∀ E ∈ sets M. m E < ∞
proof
fix E
assume E ∈ sets M
show m E < ∞
proof
show m E ≠ ∞
proof (rule ccontr)
assume ¬ m E ≠ ∞
have m E = ∞
using ⟨¬ m E ≠ ∞⟩ by auto
have signed-measure M m
using ⟨signed-measure-space M m⟩ signed-measure-space-def by auto
moreover have m E = − μ E using ⟨m = (λA . − μ A)⟩ by auto
then have ∞ ∉ range m using ⟨signed-measure M m⟩
by (metis (no-types, lifting) False ereal-less-PInfty
ereal-uminus-eq-reorder image-iff inf-range m-def rangeI)
show False using ⟨m E = ∞⟩ ⟨∞ ∉ range m⟩
by (metis rangeI)
qed
qed
qed
qed
hence ∃ M1 M2. (neg-meas-set M1) ∧ (pos-meas-set M2) ∧ (space M = M1 ∪
M2) ∧
(M1 ∩ M2 = {})
using pos-meas-set-opp neg-meas-set-opp unfolding m-def
by (metis sgn-meas signed-measure-minus signed-measure-space-def
signed-measure-space.hahn-space-decomp-def)
thus ?thesis using hahn-space-decomp-def by (metis inf-commute sup-commute)
qed

```

5 The Jordan decomposition theorem

definition *jordan-decomp* where

```
jordan-decomp m1 m2 ↔ ((measure-space (space M) (sets M) m1) ∧
  (measure-space (space M) (sets M) m2) ∧
  (∀ A ∈ sets M. 0 ≤ m1 A) ∧
  (∀ A ∈ sets M. 0 ≤ m2 A) ∧
  (∀ A ∈ sets M. μ A = (m1 A) − (m2 A)) ∧
  (∀ P N A. hahn-space-decomp P N →
    (A ∈ sets M → A ⊆ P → (m2 A) = 0) ∧
    (A ∈ sets M → A ⊆ N → (m1 A) = 0)) ∧
  ((∀ A ∈ sets M. m1 A < ∞) ∨ (∀ A ∈ sets M. m2 A < ∞)))
```

lemma *jordan-decomp-pos-meas*:

```
assumes jordan-decomp m1 m2
and hahn-space-decomp P N
and A ∈ sets M
shows m1 A = μ (A ∩ P)
```

proof –

```
have A ∩ P ∈ sets M using assms unfolding hahn-space-decomp-def
  by (simp add: pos-meas-setD1 sets.Int)
have A ∩ N ∈ sets M using assms unfolding hahn-space-decomp-def
  by (simp add: neg-meas-setD1 sets.Int)
have (A ∩ P) ∩ (A ∩ N) = {} using assms unfolding hahn-space-decomp-def
  by auto
have A = (A ∩ P) ∪ (A ∩ N) using assms unfolding hahn-space-decomp-def
  by (metis Int-Un-distrib sets.Int-space-eq2)
hence m1 A = m1 ((A ∩ P) ∪ (A ∩ N)) by simp
also have ... = m1 (A ∩ P) + m1 (A ∩ N)
  using assms pos-e2ennreal-additive[of M m1] ⟨A ∩ P ∈ sets M⟩ ⟨A ∩ N ∈ sets M⟩
  ⟨A ∩ P ∩ (A ∩ N) = {}⟩
  unfolding jordan-decomp-def additive-def by simp
also have ... = m1 (A ∩ P) using assms unfolding jordan-decomp-def
  by (metis Int-lower2 ⟨A ∩ N ∈ sets M⟩ add.right-neutral)
also have ... = m1 (A ∩ P) − m2 (A ∩ P)
  using assms unfolding jordan-decomp-def
  by (metis Int-subset-iff ⟨A ∩ P ∈ sets M⟩ ereal-minus(7)
    local.pos-wtn-base pos-wtn-subset)
also have ... = μ (A ∩ P) using assms ⟨A ∩ P ∈ sets M⟩
  unfolding jordan-decomp-def by simp
finally show ?thesis .
```

qed

lemma *jordan-decomp-neg-meas*:

```
assumes jordan-decomp m1 m2
and hahn-space-decomp P N
and A ∈ sets M
shows m2 A = −μ (A ∩ N)
```

```

proof -
have  $A \cap P \in \text{sets } M$  using assms unfolding hahn-space-decomp-def
  by (simp add: pos-meas-setD1 sets.Int)
have  $A \cap N \in \text{sets } M$  using assms unfolding hahn-space-decomp-def
  by (simp add: neg-meas-setD1 sets.Int)
have  $(A \cap P) \cap (A \cap N) = \{\}$ 
  using assms unfolding hahn-space-decomp-def by auto
have  $A = (A \cap P) \cup (A \cap N)$ 
  using assms unfolding hahn-space-decomp-def
  by (metis Int-Un-distrib sets.Int-space-eq2)
hence  $m2 A = m2 ((A \cap P) \cup (A \cap N))$  by simp
also have ... =  $m2 (A \cap P) + m2 (A \cap N)$ 
  using pos-e2ennreal-additive[of M m2] assms
  ⟨ $A \cap P \in \text{sets } M$ ⟩ ⟨ $A \cap N \in \text{sets } M$ ⟩ ⟨ $A \cap P \cap (A \cap N) = \{\}$ ⟩
  unfolding jordan-decomp-def additive-def by simp
also have ... =  $m2 (A \cap N)$  using assms unfolding jordan-decomp-def
  by (metis Int-lower2 ⟨ $A \cap P \in \text{sets } M$ ⟩ add.commute add.right-neutral)
also have ... =  $m2 (A \cap N) - m1 (A \cap N)$ 
  using assms unfolding jordan-decomp-def
  by (metis Int-lower2 ⟨ $A \cap N \in \text{sets } M$ ⟩ ereal-minus(7))
also have ... =  $-\mu (A \cap N)$  using assms ⟨ $A \cap P \in \text{sets } M$ ⟩
  unfolding jordan-decomp-def
  by (metis Diff-cancel Diff-eq-empty-iff Int-Un-eq(2) ⟨ $A \cap N \in \text{sets } M$ ⟩
    ⟨ $m2 (A \cap N) = m2 (A \cap N) - m1 (A \cap N)$ ⟩ ereal-minus(8)
    ereal-uminus-eq-reorder sup.bounded-iff)
finally show ?thesis .
qed

lemma pos-inter-neg-0:
assumes hahn-space-decomp M1 M2
  and hahn-space-decomp P N
  and  $A \in \text{sets } M$ 
  and  $A \subseteq N$ 
shows  $\mu (A \cap M1) = 0$ 
proof -
have  $\mu (A \cap M1) = \mu (A \cap ((M1 \cap P) \cup (M1 \cap (\text{sym-diff } M1 P))))$ 
  by (metis Diff-subset-conv Int-Un-distrib Un-upper1 inf.orderE)
also have ... =  $\mu ((A \cap (M1 \cap P)) \cup (A \cap (M1 \cap (\text{sym-diff } M1 P))))$ 
  by (simp add: Int-Un-distrib)
also have ... =  $\mu (A \cap (M1 \cap P)) + \mu (A \cap (M1 \cap (\text{sym-diff } M1 P)))$ 
proof (rule signed-measure-add)
  show signed-measure M  $\mu$  using sgn-meas .
  show  $A \cap (M1 \cap P) \in \text{sets } M$ 
    by (meson assms(1) assms(2) assms(3) hahn-space-decomp-def sets.Int
      signed-measure-space.pos-meas-setD1 signed-measure-space-axioms)
  show  $A \cap (M1 \cap \text{sym-diff } M1 P) \in \text{sets } M$ 
    by (meson Diff-subset assms(1) assms(2) assms(3) hahn-space-decomp-def
      pos-meas-setD1 pos-meas-set-union pos-meas-subset sets.Diff sets.Int)
  show  $A \cap (M1 \cap P) \cap (A \cap (M1 \cap \text{sym-diff } M1 P)) = \{\}$  by auto

```

```

qed
also have ... =  $\mu(A \cap (M1 \cap (\text{sym-diff } M1 P)))$ 
proof -
  have  $A \cap (M1 \cap P) = \{\}$  using assms hahn-space-decomp-def by auto
  thus ?thesis using signed-measure-empty[OF sgn-meas] by simp
qed
also have ... = 0
proof (rule hahn-decomp-ess-unique[OF assms(1) assms(2)])
  show  $A \cap (M1 \cap \text{sym-diff } M1 P) \subseteq \text{sym-diff } M1 P \cup \text{sym-diff } M2 N$  by auto
  show  $A \cap (M1 \cap \text{sym-diff } M1 P) \in \text{sets } M$ 
proof -
  have  $\text{sym-diff } M1 P \in \text{sets } M$  using assms
  by (meson hahn-space-decomp-def sets.Diff sets.Un
       signed-measure-space.pos-meas-setD1 signed-measure-space-axioms)
  hence  $M1 \cap \text{sym-diff } M1 P \in \text{sets } M$ 
  by (meson assms(1) hahn-space-decomp-def pos-meas-setD1 sets.Int)
  thus ?thesis by (simp add: assms sets.Int)
qed
qed
finally show ?thesis .
qed

lemma neg-inter-pos-0:
assumes hahn-space-decomp M1 M2
and hahn-space-decomp P N
and  $A \in \text{sets } M$ 
and  $A \subseteq P$ 
shows  $\mu(A \cap M2) = 0$ 
proof -
  have  $\mu(A \cap M2) = \mu(A \cap ((M2 \cap N) \cup (M2 \cap (\text{sym-diff } M2 N))))$ 
  by (metis Diff-subset-conv Int-Un-distrib Un-upper1 inf.orderE)
  also have ... =  $\mu((A \cap (M2 \cap N)) \cup (A \cap (M2 \cap (\text{sym-diff } M2 N))))$ 
  by (simp add: Int-Un-distrib)
  also have ... =  $\mu(A \cap (M2 \cap N)) + \mu(A \cap (M2 \cap (\text{sym-diff } M2 N)))$ 
  proof (rule signed-measure-add)
    show signed-measure M  $\mu$  using sgn-meas .
    show  $A \cap (M2 \cap N) \in \text{sets } M$ 
    by (meson assms(1) assms(2) assms(3) hahn-space-decomp-def sets.Int
         signed-measure-space.neg-meas-setD1 signed-measure-space-axioms)
    show  $A \cap (M2 \cap \text{sym-diff } M2 N) \in \text{sets } M$ 
    by (meson Diff-subset assms(1) assms(2) assms(3) hahn-space-decomp-def
         neg-meas-setD1 neg-meas-set-union neg-meas-subset sets.Diff sets.Int)
    show  $A \cap (M2 \cap N) \cap (A \cap (M2 \cap \text{sym-diff } M2 N)) = \{\}$  by auto
  qed
  also have ... =  $\mu(A \cap (M2 \cap (\text{sym-diff } M2 N)))$ 
  proof -
    have  $A \cap (M2 \cap N) = \{\}$  using assms hahn-space-decomp-def by auto
    thus ?thesis using signed-measure-empty[OF sgn-meas] by simp
  qed

```

```

also have ... = 0
proof (rule hahn-decomp-ess-unique[OF assms(1) assms(2)])
  show A ∩ (M2 ∩ sym-diff M2 N) ⊆ sym-diff M1 P ∪ sym-diff M2 N by auto
  show A ∩ (M2 ∩ sym-diff M2 N) ∈ sets M
proof -
  have sym-diff M2 N ∈ sets M using assms
  by (meson hahn-space-decomp-def sets.Diff sets.Un
      signed-measure-space.neg-meas-setD1 signed-measure-space-axioms)
  hence M2 ∩ sym-diff M2 N ∈ sets M
  by (meson assms(1) hahn-space-decomp-def neg-meas-setD1 sets.Int)
  thus ?thesis by (simp add: assms sets.Int)
qed
qed
finally show ?thesis .
qed

lemma jordan-decomposition :
shows ∃ m1 m2. jordan-decomp m1 m2
proof -
  have ∃ M1 M2. hahn-space-decomp M1 M2 using hahn-decomposition
  unfolding hahn-space-decomp-def by simp
  from this obtain M1 M2 where hahn-space-decomp M1 M2 by auto
  note Mprops = this
  define m1 where m1 = (λA. μ (A ∩ M1))
  define m2 where m2 = (λA. -μ (A ∩ M2))
  show ?thesis unfolding jordan-decomp-def
  proof (intro exI allI impI conjI ballI)
    show measure-space (space M) (sets M) (λx. e2ennreal (m1 x))
    using pos-signed-to-meas-space Mprops m1-def
    unfolding hahn-space-decomp-def by auto
  next
    show measure-space (space M) (sets M) (λx. e2ennreal (m2 x))
    using neg-signed-to-meas-space Mprops m2-def
    unfolding hahn-space-decomp-def by auto
  next
    fix A
    assume A ∈ sets M
    thus 0 ≤ m1 A unfolding m1-def using Mprops
    unfolding hahn-space-decomp-def
    by (meson inf-sup-ord(2) pos-meas-setD1 sets.Int
        signed-measure-space.pos-measure-meas signed-measure-space-axioms)
  next
    fix A
    assume A ∈ sets M
    thus 0 ≤ m2 A unfolding m2-def using Mprops
    unfolding hahn-space-decomp-def
    by (metis ereal-0-le-uminus-iff inf-sup-ord(2) neg-meas-self
        neg-meas-setD1 neg-meas-subset sets.Int)
  next

```

```

fix A
assume A ∈ sets M
have μ A = μ ((A ∩ M1) ∪ (A ∩ M2)) using Mprops
  unfolding hahn-space-decomp-def
  by (metis Int-Un-distrib ⟨A ∈ sets M⟩ sets.Int-space-eq2)
also have ... = μ (A ∩ M1) + μ (A ∩ M2)
proof (rule signed-measure-add)
  show signed-measure M μ using sgn-meas .
  show A ∩ M1 ∈ sets M using Mprops ⟨A ∈ sets M⟩
    unfolding hahn-space-decomp-def
    by (simp add: pos-meas-setD1 sets.Int)
  show A ∩ M2 ∈ sets M using Mprops ⟨A ∈ sets M⟩
    unfolding hahn-space-decomp-def
    by (simp add: neg-meas-setD1 sets.Int)
  show A ∩ M1 ∩ (A ∩ M2) = {} using Mprops
    unfolding hahn-space-decomp-def by auto
qed
also have ... = m1 A - m2 A using m1-def m2-def by simp
finally show μ A = m1 A - m2 A .
next
fix P N A
assume hahn-space-decomp P N and A ∈ sets M and A ⊆ N
note hn = this
have μ (A ∩ M1) = 0
proof (rule pos-inter-neg-0[OF - hn])
  show hahn-space-decomp M1 M2 using Mprops
    unfolding hahn-space-decomp-def by simp
qed
thus m1 A = 0 unfolding m1-def by simp
next
fix P N A
assume hahn-space-decomp P N and A ∈ sets M and A ⊆ P
note hp = this
have μ (A ∩ M2) = 0
proof (rule neg-inter-pos-0[OF - hp])
  show hahn-space-decomp M1 M2 using Mprops
    unfolding hahn-space-decomp-def by simp
qed
thus m2 A = 0 unfolding m2-def by simp
next
show (∀ E ∈ sets M. m1 E < ∞) ∨ (∀ E ∈ sets M. m2 E < ∞)
proof (cases ∀ E ∈ sets M. m1 E < ∞)
  case True
  thus ?thesis by simp
next
case False
have ∀ E ∈ sets M. m2 E < ∞
proof
  fix E

```

```

assume  $E \in \text{sets } M$ 
show  $m2 E < \infty$ 
proof -
  have  $(m2 E) = -\mu(E \cap M2)$  using m2-def by simp
  also have ...  $\neq \infty$  using False sgn-meas inf-range
    by (metis ereal-less-PInfty ereal-uminus-uminus m1-def rangeI)
  finally have  $m2 E \neq \infty$  .
  thus ?thesis by (simp add: top.not-eq-extremum)
  qed
  qed
  thus ?thesis by simp
  qed
  qed
  qed

lemma jordan-decomposition-unique :
  assumes jordan-decomp m1 m2
  and jordan-decomp n1 n2
  and  $A \in \text{sets } M$ 
  shows  $m1 A = n1 A$   $m2 A = n2 A$ 
proof -
  have  $\exists M1 M2. \text{hahn-space-decomp } M1 M2$  using hahn-decomposition by simp
  with jordan-decomp-neg-meas jordan-decomp-pos-meas
  show  $m1 A = n1 A$   $m2 A = n2 A$  by (metis assms)+
  qed

end

end

```

References

- [1] E. DiBenedetto. *Real Analysis*. Birkhäuser Advanced Texts. Birkhäuser.