

Gröbner Bases Theory

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Abstract

This formalization is concerned with the theory of Gröbner bases in (commutative) multivariate polynomial rings over fields, originally developed by Buchberger in his 1965 PhD thesis. Apart from the statement and proof of the main theorem of the theory, the formalization also implements algorithms for actually computing Gröbner bases, thus allowing to effectively decide ideal membership in finitely generated polynomial ideals. Furthermore, all functions can be executed on a concrete representation of multivariate polynomials as association lists.

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1 Introduction

The theory of Gröbner bases, invented by Buchberger in [2, 3], is ubiquitous in many areas of computer algebra and beyond, as it allows to effectively solve a multitude of interesting, non-trivial problems of polynomial ideal theory. Since its invention in the mid-sixties, the theory has already seen a whole range of extensions and generalizations, some of which are present in this formalization:

- Following [11], the theory is formulated for vector-polynomials instead of ordinary scalar polynomials, thus allowing to compute Gröbner bases of syzygy modules.
- Besides Buchberger’s original algorithm, the formalization also features Faugère’s F_4 algorithm [8] for computing Gröbner bases.
- All algorithms for computing Gröbner bases incorporate criteria to avoid useless pairs; see [4] for details.
- Reduced Gröbner bases have been formalized and can be computed by a formally verified algorithm, too.

For further information about Gröbner bases theory the interested reader may consult the introductory paper [5] or literally any book on commutative/computer algebra, e. g. [1, 11].

1.1 Related Work

The theory of Gröbner bases has already been formalized in a couple of other proof assistants, listed below in alphabetical order:

- ACL2 [13],
- Coq [16, 10],
- Mizar [15], and
- Theorema [6, 12].

Please note that this formalization must not be confused with the *algebra* proof method based on Gröbner bases [7], which is a completely independent piece of work: our results could in principle be used to formally prove the correctness and, to some extent, completeness of said proof method.

1.2 Future Work

This formalization can be extended in several ways:

- One could formalize signature-based algorithms for computing Gröbner bases, as for instance Faugère's F_5 algorithm [9]. Such algorithms are typically more efficient than Buchberger's algorithm.
- One could establish the connection to *elimination theory*, exploiting the well-known *elimination property* of Gröbner bases w. r. t. certain term-orders (e. g. the purely lexicographic one). This would enable the effective simplification (and even solution, in some sense) of systems of algebraic equations.
- One could generalize the theory further to cover also *non-commutative* Gröbner bases [14].

2 General Utilities

```
theory General
  imports Polynomials.Utils
begin
```

A couple of general-purpose functions and lemmas, mainly related to lists.

2.1 Lists

```
lemma distinct-reorder: distinct (xs @ (y # ys)) = distinct (y # (xs @ ys)) <proof>
```

```
lemma set-reorder: set (xs @ (y # ys)) = set (y # (xs @ ys)) <proof>
```

```
lemma distinctI:
```

```
  assumes  $\bigwedge i j. i < j \implies i < \text{length } xs \implies j < \text{length } xs \implies xs ! i \neq xs ! j$ 
  shows distinct xs
  <proof>
```

```
lemma filter-nth-pairE:
```

```
  assumes  $i < j$  and  $i < \text{length } (\text{filter } P \text{ } xs)$  and  $j < \text{length } (\text{filter } P \text{ } xs)$ 
  obtains  $i' j'$  where  $i' < j'$  and  $i' < \text{length } xs$  and  $j' < \text{length } xs$ 
    and  $(\text{filter } P \text{ } xs) ! i = xs ! i'$  and  $(\text{filter } P \text{ } xs) ! j = xs ! j'$ 
  <proof>
```

```
lemma distinct-filterI:
```

```
  assumes  $\bigwedge i j. i < j \implies i < \text{length } xs \implies j < \text{length } xs \implies P (xs ! i) \implies P (xs ! j) \implies xs ! i \neq xs ! j$ 
  shows distinct (filter P xs)
  <proof>
```

lemma *set- zip -map*: $\text{set} (\text{zip} (\text{map } f \text{ } xs) (\text{map } g \text{ } xs)) = (\lambda x. (f \ x, g \ x)) \text{ ` } (\text{set } xs)$
 <proof>

lemma *set- zip -map1*: $\text{set} (\text{zip} (\text{map } f \text{ } xs) \text{ } xs) = (\lambda x. (f \ x, x)) \text{ ` } (\text{set } xs)$
 <proof>

lemma *set- zip -map2*: $\text{set} (\text{zip } xs (\text{map } f \text{ } xs)) = (\lambda x. (x, f \ x)) \text{ ` } (\text{set } xs)$
 <proof>

lemma *UN-upt*: $(\bigcup i \in \{0..<\text{length } xs\}. f \ (xs \ ! \ i)) = (\bigcup x \in \text{set } xs. f \ x)$
 <proof>

lemma *sum-list-zeroI'*:
 assumes $\bigwedge i. i < \text{length } xs \implies xs \ ! \ i = 0$
 shows $\text{sum-list } xs = 0$
 <proof>

lemma *sum-list-map2-plus*:
 assumes $\text{length } xs = \text{length } ys$
 shows $\text{sum-list} (\text{map2 } (+) \text{ } xs \ ys) = \text{sum-list } xs + \text{sum-list} (ys :: 'a :: \text{comm-monoid-add list})$
 <proof>

lemma *sum-list-eq-nthI*:
 assumes $i < \text{length } xs$ and $\bigwedge j. j < \text{length } xs \implies j \neq i \implies xs \ ! \ j = 0$
 shows $\text{sum-list } xs = xs \ ! \ i$
 <proof>

2.1.1 *max-list*

fun (in *ord*) *max-list* :: 'a list \Rightarrow 'a **where**
 $\text{max-list } (x \ \# \ xs) = (\text{case } xs \ \text{of } [] \Rightarrow x \ | \ - \Rightarrow \text{max } x \ (\text{max-list } xs))$

context *linorder*
begin

lemma *max-list-Max*: $xs \neq [] \implies \text{max-list } xs = \text{Max} (\text{set } xs)$
 <proof>

lemma *max-list-ge*:
 assumes $x \in \text{set } xs$
 shows $x \leq \text{max-list } xs$
 <proof>

lemma *max-list-boundedI*:
 assumes $xs \neq []$ and $\bigwedge x. x \in \text{set } xs \implies x \leq a$
 shows $\text{max-list } xs \leq a$
 <proof>

end

2.1.2 *insort-wrt*

primrec *insort-wrt* :: ('c ⇒ 'c ⇒ bool) ⇒ 'c ⇒ 'c list ⇒ 'c list **where**
 insort-wrt - x [] = [x] |
 insort-wrt r x (y # ys) =
 (if r x y then (x # y # ys) else y # (*insort-wrt* r x ys))

lemma *insort-wrt-not-Nil* [simp]: *insort-wrt* r x xs ≠ []
 ⟨proof⟩

lemma *length-insort-wrt* [simp]: *length* (*insort-wrt* r x xs) = *Suc* (*length* xs)
 ⟨proof⟩

lemma *set-insort-wrt* [simp]: *set* (*insort-wrt* r x xs) = *insert* x (*set* xs)
 ⟨proof⟩

lemma *sorted-wrt-insort-wrt-imp-sorted-wrt*:
 assumes *sorted-wrt* r (*insort-wrt* s x xs)
 shows *sorted-wrt* r xs
 ⟨proof⟩

lemma *sorted-wrt-imp-sorted-wrt-insort-wrt*:
 assumes *transp* r **and** $\bigwedge a. r a x \vee r x a$ **and** *sorted-wrt* r xs
 shows *sorted-wrt* r (*insort-wrt* r x xs)
 ⟨proof⟩

corollary *sorted-wrt-insort-wrt*:
 assumes *transp* r **and** $\bigwedge a. r a x \vee r x a$
 shows *sorted-wrt* r (*insort-wrt* r x xs) \longleftrightarrow *sorted-wrt* r xs (**is** ?l \longleftrightarrow ?r)
 ⟨proof⟩

2.1.3 *diff-list* **and** *insert-list*

definition *diff-list* :: 'a list ⇒ 'a list ⇒ 'a list (**infixl** -- 65)
 where *diff-list* xs ys = *fold* *removeAll* ys xs

lemma *set-diff-list*: *set* (xs -- ys) = *set* xs - *set* ys
 ⟨proof⟩

lemma *diff-list-disjoint*: *set* ys ∩ *set* (xs -- ys) = {}
 ⟨proof⟩

lemma *subset-append-diff-cancel*:
 assumes *set* ys ⊆ *set* xs
 shows *set* (ys @ (xs -- ys)) = *set* xs
 ⟨proof⟩

definition *insert-list* :: 'a ⇒ 'a list ⇒ 'a list

where $insert\text{-}list\ x\ xs = (if\ x \in set\ xs\ then\ xs\ else\ x\ \# \ xs)$

lemma $set\text{-}insert\text{-}list: set\ (insert\text{-}list\ x\ xs) = insert\ x\ (set\ xs)$
(*proof*)

2.1.4 $remdups\text{-}wrt$

primrec $remdups\text{-}wrt :: ('a \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow 'a\ list\ \mathbf{where}$
 $remdups\text{-}wrt\text{-}base: remdups\text{-}wrt\ -\ [] = []\ |$
 $remdups\text{-}wrt\text{-}rec: remdups\text{-}wrt\ f\ (x\ \# \ xs) = (if\ f\ x \in f\ ' \ set\ xs\ then\ remdups\text{-}wrt\ f\ xs\ else\ x\ \# \ remdups\text{-}wrt\ f\ xs)$

lemma $set\text{-}remdups\text{-}wrt: f\ ' \ set\ (remdups\text{-}wrt\ f\ xs) = f\ ' \ set\ xs$
(*proof*)

lemma $subset\text{-}remdups\text{-}wrt: set\ (remdups\text{-}wrt\ f\ xs) \subseteq set\ xs$
(*proof*)

lemma $remdups\text{-}wrt\text{-}distinct\text{-}wrt:$
assumes $x \in set\ (remdups\text{-}wrt\ f\ xs)$ **and** $y \in set\ (remdups\text{-}wrt\ f\ xs)$ **and** $x \neq y$
shows $f\ x \neq f\ y$
(*proof*)

lemma $distinct\text{-}remdups\text{-}wrt: distinct\ (remdups\text{-}wrt\ f\ xs)$
(*proof*)

lemma $map\text{-}remdups\text{-}wrt: map\ f\ (remdups\text{-}wrt\ f\ xs) = remdups\ (map\ f\ xs)$
(*proof*)

lemma $remdups\text{-}wrt\text{-}append:$
 $remdups\text{-}wrt\ f\ (xs\ @ \ ys) = (filter\ (\lambda a. f\ a \notin f\ ' \ set\ ys)\ (remdups\text{-}wrt\ f\ xs))\ @$
 $(remdups\text{-}wrt\ f\ ys)$
(*proof*)

2.1.5 $map\text{-}idx$

primrec $map\text{-}idx :: ('a \Rightarrow nat \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow nat \Rightarrow 'b\ list\ \mathbf{where}$
 $map\text{-}idx\ f\ []\ n = []\ |$
 $map\text{-}idx\ f\ (x\ \# \ xs)\ n = (f\ x\ n)\ \# \ (map\text{-}idx\ f\ xs\ (Suc\ n))$

lemma $map\text{-}idx\text{-}eq\text{-}map2: map\text{-}idx\ f\ xs\ n = map2\ f\ xs\ [n..<n + length\ xs]$
(*proof*)

lemma $length\text{-}map\text{-}idx\ [simp]: length\ (map\text{-}idx\ f\ xs\ n) = length\ xs$
(*proof*)

lemma $map\text{-}idx\text{-}append: map\text{-}idx\ f\ (xs\ @ \ ys)\ n = (map\text{-}idx\ f\ xs\ n)\ @ \ (map\text{-}idx\ f\ ys\ (n + length\ xs))$
(*proof*)

lemma *map-idx-nth*:

assumes $i < \text{length } xs$

shows $(\text{map-idx } f \text{ } xs \ n) \ ! \ i = f \ (xs \ ! \ i) \ (n + i)$

<proof>

lemma *map-map-idx*: $\text{map } f \ (\text{map-idx } g \ xs \ n) = \text{map-idx } (\lambda x \ i. f \ (g \ x \ i)) \ xs \ n$

<proof>

lemma *map-idx-map*: $\text{map-idx } f \ (\text{map } g \ xs) \ n = \text{map-idx } (f \circ g) \ xs \ n$

<proof>

lemma *map-idx-no-idx*: $\text{map-idx } (\lambda x \ -. \ f \ x) \ xs \ n = \text{map } f \ xs$

<proof>

lemma *map-idx-no-elem*: $\text{map-idx } (\lambda \cdot. f) \ xs \ n = \text{map } f \ [n..<n + \text{length } xs]$

<proof>

lemma *map-idx-eq-map*: $\text{map-idx } f \ xs \ n = \text{map } (\lambda i. f \ (xs \ ! \ i) \ (i + n)) \ [0..<\text{length } xs]$

<proof>

lemma *set-map-idx*: $\text{set } (\text{map-idx } f \ xs \ n) = (\lambda i. f \ (xs \ ! \ i) \ (i + n)) \ ` \ \{0..<\text{length } xs\}$

<proof>

2.1.6 *map-dup*

primrec *map-dup* :: $(a \Rightarrow b) \Rightarrow (a \Rightarrow b) \Rightarrow 'a \ \text{list} \Rightarrow 'b \ \text{list}$ **where**

map-dup - - $\ [] = []$

map-dup $f \ g \ (x \ \# \ xs) = (\text{if } x \in \text{set } xs \ \text{then } g \ x \ \text{else } f \ x) \ \# \ (\text{map-dup } f \ g \ xs)$

lemma *length-map-dup[simp]*: $\text{length } (\text{map-dup } f \ g \ xs) = \text{length } xs$

<proof>

lemma *map-dup-distinct*:

assumes *distinct* xs

shows $\text{map-dup } f \ g \ xs = \text{map } f \ xs$

<proof>

lemma *filter-map-dup-const*:

$\text{filter } (\lambda x. x \neq c) \ (\text{map-dup } f \ (\lambda \cdot. c) \ xs) = \text{filter } (\lambda x. x \neq c) \ (\text{map } f \ (\text{remdups } xs))$

<proof>

lemma *filter-zip-map-dup-const*:

$\text{filter } (\lambda(a, b). a \neq c) \ (\text{zip } (\text{map-dup } f \ (\lambda \cdot. c) \ xs) \ xs) =$

$\text{filter } (\lambda(a, b). a \neq c) \ (\text{zip } (\text{map } f \ (\text{remdups } xs)) \ (\text{remdups } xs))$

<proof>

2.1.7 Filtering Minimal Elements

context

fixes $rel :: 'a \Rightarrow 'a \Rightarrow bool$

begin

primrec $filter\text{-}min\text{-}aux :: 'a\ list \Rightarrow 'a\ list \Rightarrow 'a\ list$ **where**

$filter\text{-}min\text{-}aux []\ ys = ys$

$filter\text{-}min\text{-}aux (x \# xs)\ ys =$

(if $(\exists y \in (set\ xs \cup set\ ys). rel\ y\ x)$ then $(filter\text{-}min\text{-}aux\ xs\ ys)$

else $(filter\text{-}min\text{-}aux\ xs\ (x \# ys))$)

definition $filter\text{-}min :: 'a\ list \Rightarrow 'a\ list$

where $filter\text{-}min\ xs = filter\text{-}min\text{-}aux\ xs\ []$

definition $filter\text{-}min\text{-}append :: 'a\ list \Rightarrow 'a\ list \Rightarrow 'a\ list$

where $filter\text{-}min\text{-}append\ xs\ ys =$

(let $P = (\lambda zs. \lambda x. \neg (\exists z \in set\ zs. rel\ z\ x))$; $ys1 = filter\ (P\ xs)\ ys$ in

$(filter\ (P\ ys1)\ xs) @ ys1$)

lemma $filter\text{-}min\text{-}aux\text{-}supset: set\ ys \subseteq set\ (filter\text{-}min\text{-}aux\ xs\ ys)$

<proof>

lemma $filter\text{-}min\text{-}aux\text{-}subset: set\ (filter\text{-}min\text{-}aux\ xs\ ys) \subseteq set\ xs \cup set\ ys$

<proof>

lemma $filter\text{-}min\text{-}aux\text{-}relE:$

assumes $transp\ rel$ **and** $x \in set\ xs$ **and** $x \notin set\ (filter\text{-}min\text{-}aux\ xs\ ys)$

obtains y **where** $y \in set\ (filter\text{-}min\text{-}aux\ xs\ ys)$ **and** $rel\ y\ x$

<proof>

lemma $filter\text{-}min\text{-}aux\text{-}minimal:$

assumes $transp\ rel$ **and** $x \in set\ (filter\text{-}min\text{-}aux\ xs\ ys)$ **and** $y \in set\ (filter\text{-}min\text{-}aux\ xs\ ys)$

and $rel\ x\ y$

assumes $\bigwedge a\ b. a \in set\ xs \cup set\ ys \implies b \in set\ ys \implies rel\ a\ b \implies a = b$

shows $x = y$

<proof>

lemma $filter\text{-}min\text{-}aux\text{-}distinct:$

assumes $reflp\ rel$ **and** $distinct\ ys$

shows $distinct\ (filter\text{-}min\text{-}aux\ xs\ ys)$

<proof>

lemma $filter\text{-}min\text{-}subset: set\ (filter\text{-}min\ xs) \subseteq set\ xs$

<proof>

lemma $filter\text{-}min\text{-}cases:$

assumes $transp\ rel$ **and** $x \in set\ xs$

assumes $x \in set\ (filter\text{-}min\ xs) \implies thesis$

assumes $\bigwedge y. y \in \text{set } (\text{filter-min } xs) \implies x \notin \text{set } (\text{filter-min } xs) \implies \text{rel } y \ x \implies$
thesis
shows *thesis*
 ⟨*proof*⟩

corollary *filter-min-relE*:

assumes *transp rel* **and** *reflp rel* **and** $x \in \text{set } xs$
obtains y **where** $y \in \text{set } (\text{filter-min } xs)$ **and** $\text{rel } y \ x$
 ⟨*proof*⟩

lemma *filter-min-minimal*:

assumes *transp rel* **and** $x \in \text{set } (\text{filter-min } xs)$ **and** $y \in \text{set } (\text{filter-min } xs)$ **and**
rel x y
shows $x = y$
 ⟨*proof*⟩

lemma *filter-min-distinct*:

assumes *reflp rel*
shows *distinct (filter-min xs)*
 ⟨*proof*⟩

lemma *filter-min-append-subset*: $\text{set } (\text{filter-min-append } xs \ ys) \subseteq \text{set } xs \cup \text{set } ys$
 ⟨*proof*⟩

lemma *filter-min-append-cases*:

assumes *transp rel* **and** $x \in \text{set } xs \cup \text{set } ys$
assumes $x \in \text{set } (\text{filter-min-append } xs \ ys) \implies$ *thesis*
assumes $\bigwedge y. y \in \text{set } (\text{filter-min-append } xs \ ys) \implies x \notin \text{set } (\text{filter-min-append } xs$
ys) \implies rel y x \implies thesis
shows *thesis*
 ⟨*proof*⟩

corollary *filter-min-append-relE*:

assumes *transp rel* **and** *reflp rel* **and** $x \in \text{set } xs \cup \text{set } ys$
obtains y **where** $y \in \text{set } (\text{filter-min-append } xs \ ys)$ **and** $\text{rel } y \ x$
 ⟨*proof*⟩

lemma *filter-min-append-minimal*:

assumes $\bigwedge x' \ y'. x' \in \text{set } xs \implies y' \in \text{set } xs \implies \text{rel } x' \ y' \implies x' = y'$
and $\bigwedge x' \ y'. x' \in \text{set } ys \implies y' \in \text{set } ys \implies \text{rel } x' \ y' \implies x' = y'$
and $x \in \text{set } (\text{filter-min-append } xs \ ys)$ **and** $y \in \text{set } (\text{filter-min-append } xs \ ys)$
and *rel x y*
shows $x = y$
 ⟨*proof*⟩

lemma *filter-min-append-distinct*:

assumes *reflp rel* **and** *distinct xs* **and** *distinct ys*
shows *distinct (filter-min-append xs ys)*
 ⟨*proof*⟩

end

end

3 Properties of Binary Relations

theory *Confluence*

imports *Abstract-Rewriting.Abstract-Rewriting Open-Induction.Restricted-Predicates*
begin

This theory formalizes some general properties of binary relations, in particular a very weak sufficient condition for a relation to be Church-Rosser.

3.1 *Restricted-Predicates.wfp-on*

lemma *wfp-on-imp-wfP*:

assumes *wfp-on r A*

shows $\text{wfP } (\lambda x y. r x y \wedge x \in A \wedge y \in A)$ (**is** *wfP ?r*)

<proof>

lemma *wfp-onI-min*:

assumes $\bigwedge x Q. x \in Q \implies Q \subseteq A \implies \exists z \in Q. \forall y \in A. r y z \longrightarrow y \notin Q$

shows *wfp-on r A*

<proof>

lemma *wfp-onE-min*:

assumes *wfp-on r A* **and** $x \in Q$ **and** $Q \subseteq A$

obtains z **where** $z \in Q$ **and** $\bigwedge y. r y z \implies y \notin Q$

<proof>

lemma *wfp-onI-chain*: $\neg (\exists f. \forall i. f i \in A \wedge r (f (Suc i)) (f i)) \implies \text{wfp-on } r A$

<proof>

lemma *finite-minimalE*:

assumes *finite A* **and** $A \neq \{\}$ **and** *irreflp rel* **and** *transp rel*

obtains a **where** $a \in A$ **and** $\bigwedge b. rel b a \implies b \notin A$

<proof>

lemma *wfp-on-finite*:

assumes *irreflp rel* **and** *transp rel* **and** *finite A*

shows *wfp-on rel A*

<proof>

3.2 Relations

locale *relation* = **fixes** $r::'a \Rightarrow 'a \Rightarrow \text{bool}$ (**infixl** $\rightarrow 50$)

begin

abbreviation $rtc::'a \Rightarrow 'a \Rightarrow bool$ (**infixl** \rightarrow^* 50)
where $rtc\ a\ b \equiv r^{**}\ a\ b$

abbreviation $sc::'a \Rightarrow 'a \Rightarrow bool$ (**infixl** \leftrightarrow 50)
where $sc\ a\ b \equiv a \rightarrow b \vee b \rightarrow a$

definition $is-final::'a \Rightarrow bool$ **where**
 $is-final\ a \equiv \neg (\exists b. r\ a\ b)$

definition $srtc::'a \Rightarrow 'a \Rightarrow bool$ (**infixl** \leftrightarrow^* 50) **where**
 $srtc\ a\ b \equiv sc^{**}\ a\ b$

definition $cs::'a \Rightarrow 'a \Rightarrow bool$ (**infixl** \downarrow^* 50) **where**
 $cs\ a\ b \equiv (\exists s. (a \rightarrow^* s) \wedge (b \rightarrow^* s))$

definition $is-confluent-on :: 'a\ set \Rightarrow bool$
where $is-confluent-on\ A \longleftrightarrow (\forall a \in A. \forall b1\ b2. (a \rightarrow^* b1 \wedge a \rightarrow^* b2) \longrightarrow b1 \downarrow^* b2)$

definition $is-confluent :: bool$
where $is-confluent \equiv is-confluent-on\ UNIV$

definition $is-loc-confluent :: bool$
where $is-loc-confluent \equiv (\forall a\ b1\ b2. (a \rightarrow b1 \wedge a \rightarrow b2) \longrightarrow b1 \downarrow^* b2)$

definition $is-ChurchRosser :: bool$
where $is-ChurchRosser \equiv (\forall a\ b. a \leftrightarrow^* b \longrightarrow a \downarrow^* b)$

definition $dw-closed :: 'a\ set \Rightarrow bool$
where $dw-closed\ A \longleftrightarrow (\forall a \in A. \forall b. a \rightarrow b \longrightarrow b \in A)$

lemma $dw-closedI$ [*intro*]:
assumes $\bigwedge a\ b. a \in A \Longrightarrow a \rightarrow b \Longrightarrow b \in A$
shows $dw-closed\ A$
<proof>

lemma $dw-closedD$:
assumes $dw-closed\ A$ **and** $a \in A$ **and** $a \rightarrow b$
shows $b \in A$
<proof>

lemma $dw-closed-rtrancl$:
assumes $dw-closed\ A$ **and** $a \in A$ **and** $a \rightarrow^* b$
shows $b \in A$
<proof>

lemma $dw-closed-empty$: $dw-closed\ \{\}$
<proof>

lemma $dw-closed-UNIV$: $dw-closed\ UNIV$

$\langle proof \rangle$

3.3 Setup for Connection to Theory *Abstract-Rewriting.Abstract-Rewriting*

abbreviation (*input*) $relset::('a * 'a)$ set where

$relset \equiv \{(x, y). x \rightarrow y\}$

lemma *rtc-rtranclI*:

assumes $a \rightarrow^* b$

shows $(a, b) \in relset^*$

$\langle proof \rangle$

lemma *final-NF*: $(is-final\ a) = (a \in NF\ relset)$

$\langle proof \rangle$

lemma *sc-symcl*: $(a \leftrightarrow b) = ((a, b) \in relset^{\leftrightarrow})$

$\langle proof \rangle$

lemma *srtc-conversion*: $(a \leftrightarrow^* b) = ((a, b) \in relset^{\leftrightarrow*})$

$\langle proof \rangle$

lemma *cs-join*: $(a \downarrow^* b) = ((a, b) \in relset^{\downarrow})$

$\langle proof \rangle$

lemma *confluent-CR*: $is-confluent = CR\ relset$

$\langle proof \rangle$

lemma *ChurchRosser-conversion*: $is-ChurchRosser = (relset^{\leftrightarrow*} \subseteq relset^{\downarrow})$

$\langle proof \rangle$

lemma *loc-confluent-WCR*:

shows $is-loc-confluent = WCR\ relset$

$\langle proof \rangle$

lemma *wf-converse*:

shows $(wfP\ r^{\hat{-}-1}) = (wf\ (relset^{-1}))$

$\langle proof \rangle$

lemma *wf-SN*:

shows $(wfP\ r^{\hat{-}-1}) = (SN\ relset)$

$\langle proof \rangle$

3.4 Simple Lemmas

lemma *rtrancl-is-final*:

assumes $a \rightarrow^* b$ and *is-final* a

shows $a = b$

$\langle proof \rangle$

lemma *cs-refl*:

shows $x \downarrow^* x$
<proof>

lemma *cs-sym*:
assumes $x \downarrow^* y$
shows $y \downarrow^* x$
<proof>

lemma *rtc-implies-cs*:
assumes $x \rightarrow^* y$
shows $x \downarrow^* y$
<proof>

lemma *rtc-implies-srtc*:
assumes $a \rightarrow^* b$
shows $a \leftrightarrow^* b$
<proof>

lemma *srtc-symmetric*:
assumes $a \leftrightarrow^* b$
shows $b \leftrightarrow^* a$
<proof>

lemma *srtc-transitive*:
assumes $a \leftrightarrow^* b$ **and** $b \leftrightarrow^* c$
shows $a \leftrightarrow^* c$
<proof>

lemma *cs-implies-srtc*:
assumes $a \downarrow^* b$
shows $a \leftrightarrow^* b$
<proof>

lemma *confluence-equiv-ChurchRosser*: $is\text{-confluent} = is\text{-ChurchRosser}$
<proof>

corollary *confluence-implies-ChurchRosser*:
assumes $is\text{-confluent}$
shows $is\text{-ChurchRosser}$
<proof>

lemma *ChurchRosser-unique-final*:
assumes $is\text{-ChurchRosser}$ **and** $a \rightarrow^* b1$ **and** $a \rightarrow^* b2$ **and** $is\text{-final } b1$ **and**
 $is\text{-final } b2$
shows $b1 = b2$
<proof>

lemma *wf-on-imp-nf-ex*:
assumes $wfp\text{-on } ((\rightarrow)^{-1-1}) A$ **and** $dw\text{-closed } A$ **and** $a \in A$

obtains b where $a \rightarrow^* b$ and *is-final* b
 $\langle proof \rangle$

lemma *unique-nf-imp-confluence-on*:

assumes *major*: $\bigwedge a b1 b2. a \in A \implies (a \rightarrow^* b1) \implies (a \rightarrow^* b2) \implies \textit{is-final } b1$
 $\implies \textit{is-final } b2 \implies b1 = b2$
and *wf*: *wfp-on* $((\rightarrow)^{-1-1}) A$ **and *dw*:** *dw-closed* A
shows *is-confluent-on* A
 $\langle proof \rangle$

corollary *wf-imp-nf-ex*:

assumes *wfP* $((\rightarrow)^{-1-1})$
obtains b where $a \rightarrow^* b$ and *is-final* b
 $\langle proof \rangle$

corollary *unique-nf-imp-confluence*:

assumes $\bigwedge a b1 b2. (a \rightarrow^* b1) \implies (a \rightarrow^* b2) \implies \textit{is-final } b1 \implies \textit{is-final } b2$
 $\implies b1 = b2$
and *wfP* $((\rightarrow)^{-1-1})$
shows *is-confluent*
 $\langle proof \rangle$

end

3.5 Advanced Results and the Generalized Newman Lemma

definition *relbelow-on*: $'a \textit{ set} \implies ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow 'a \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool})$

where *relbelow-on* $A \textit{ ord } z \textit{ rel } a b \equiv (a \in A \wedge b \in A \wedge \textit{rel } a b \wedge \textit{ord } a z \wedge \textit{ord } b z)$

definition *cbelow-on-1*: $'a \textit{ set} \implies ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow 'a \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool})$

where *cbelow-on-1* $A \textit{ ord } z \textit{ rel} \equiv (\textit{relbelow-on } A \textit{ ord } z \textit{ rel})^{++}$

definition *cbelow-on*: $'a \textit{ set} \implies ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow 'a \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool})$

where *cbelow-on* $A \textit{ ord } z \textit{ rel } a b \equiv (a = b \wedge b \in A \wedge \textit{ord } b z) \vee \textit{cbelow-on-1 } A \textit{ ord } z \textit{ rel } a b$

Note that *cbelow-on* cannot be defined as the reflexive-transitive closure of *relbelow-on*, since it is in general not reflexive!

definition *is-loc-connective-on*: $'a \textit{ set} \implies ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow ('a \Rightarrow 'a \Rightarrow \textit{bool}) \Rightarrow \textit{bool}$

where *is-loc-connective-on* $A \textit{ ord } r \iff (\forall a \in A. \forall b1 b2. r a b1 \wedge r a b2 \longrightarrow \textit{cbelow-on } A \textit{ ord } a (\textit{relation.sc } r) b1 b2)$

Note that *Restricted-Predicates.wfp-on* is *not* the same as *SN-on*, since in the definition of *SN-on* only the *first* element of the chain must be in the

set.

lemma *cbelow-on-first-below*:

assumes *cbelow-on A ord z rel a b*

shows *ord a z*

<proof>

lemma *cbelow-on-second-below*:

assumes *cbelow-on A ord z rel a b*

shows *ord b z*

<proof>

lemma *cbelow-on-first-in*:

assumes *cbelow-on A ord z rel a b*

shows $a \in A$

<proof>

lemma *cbelow-on-second-in*:

assumes *cbelow-on A ord z rel a b*

shows $b \in A$

<proof>

lemma *cbelow-on-intro [intro]*:

assumes *main: cbelow-on A ord z rel a b and $c \in A$ and rel b c and ord c z*

shows *cbelow-on A ord z rel a c*

<proof>

lemma *cbelow-on-induct [consumes 1, case-names base step]*:

assumes *a: cbelow-on A ord z rel a b*

and *base: $a \in A \implies ord a z \implies P a$*

and *ind: $\bigwedge b c. [\![\ cbelow-on A ord z rel a b; rel b c; c \in A; ord c z; P b \!\!] \implies$*

P c

shows *P b*

<proof>

lemma *cbelow-on-symmetric*:

assumes *main: cbelow-on A ord z rel a b and symp rel*

shows *cbelow-on A ord z rel b a*

<proof>

lemma *cbelow-on-transitive*:

assumes *cbelow-on A ord z rel a b and cbelow-on A ord z rel b c*

shows *cbelow-on A ord z rel a c*

<proof>

lemma *cbelow-on-mono*:

assumes *cbelow-on A ord z rel a b and $A \subseteq B$*

shows *cbelow-on B ord z rel a b*

<proof>

```

locale relation-order = relation +
  fixes ord::'a ⇒ 'a ⇒ bool
  fixes A::'a set
  assumes trans: ord x y ⇒ ord y z ⇒ ord x z
  assumes wf: wfp-on ord A
  assumes refines:  $(\rightarrow) \leq \text{ord}^{-1-1}$ 
begin

lemma relation-refines:
  assumes a → b
  shows ord b a
  ⟨proof⟩

lemma relation-wf: wfp-on  $(\rightarrow)^{-1-1}$  A
  ⟨proof⟩

lemma rtc-implies-cbelow-on:
  assumes dw-closed A and main: a →* b and a ∈ A and ord a c
  shows cbelow-on A ord c (↔) a b
  ⟨proof⟩

lemma cs-implies-cbelow-on:
  assumes dw-closed A and a ↓* b and a ∈ A and b ∈ A and ord a c and ord b
  c
  shows cbelow-on A ord c (↔) a b
  ⟨proof⟩

The generalized Newman lemma, taken from [17]:

lemma loc-connectivity-implies-confluence:
  assumes is-loc-connective-on A ord (→) and dw-closed A
  shows is-confluent-on A
  ⟨proof⟩

end

theorem loc-connectivity-equiv-ChurchRosser:
  assumes relation-order r ord UNIV
  shows relation.is-ChurchRosser r = is-loc-connective-on UNIV ord r
  ⟨proof⟩

end

```

4 Polynomial Reduction

```

theory Reduction
imports Polynomials.MPoly-Type-Class-Ordered Confluence
begin

```

This theory formalizes the concept of *reduction* of polynomials by polyno-

mials.

context *ordered-term*

begin

definition *red-single* :: ('t \Rightarrow_0 'b::field) \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow 'a \Rightarrow bool
where *red-single* p q f t \longleftrightarrow (f \neq 0 \wedge lookup p (t \oplus lt f) \neq 0 \wedge
q = p - monom-mult ((lookup p (t \oplus lt f)) / lc f) t f)

definition *red* :: ('t \Rightarrow_0 'b::field) set \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow bool
where *red* F p q \longleftrightarrow ($\exists f \in F. \exists t. \text{red-single } p \ q \ f \ t$)

definition *is-red* :: ('t \Rightarrow_0 'b::field) set \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow bool
where *is-red* F a \longleftrightarrow \neg relation.is-final (red F) a

4.1 Basic Properties of Reduction

lemma *red-setI*:

assumes f \in F **and** a: *red-single* p q f t
shows *red* F p q
<proof>

lemma *red-setE*:

assumes *red* F p q
obtains f **and** t **where** f \in F **and** *red-single* p q f t
<proof>

lemma *red-empty*: \neg *red* {} p q
<proof>

lemma *red-singleton-zero*: \neg *red* {0} p q
<proof>

lemma *red-union*: *red* (F \cup G) p q = (*red* F p q \vee *red* G p q)
<proof>

lemma *red-unionI1*:

assumes *red* F p q
shows *red* (F \cup G) p q
<proof>

lemma *red-unionI2*:

assumes *red* G p q
shows *red* (F \cup G) p q
<proof>

lemma *red-subset*:

assumes *red* G p q **and** G \subseteq F
shows *red* F p q
<proof>

lemma *red-union-singleton-zero*: $\text{red } (F \cup \{0\}) = \text{red } F$
<proof>

lemma *red-minus-singleton-zero*: $\text{red } (F - \{0\}) = \text{red } F$
<proof>

lemma *red-rtrancl-subset*:
assumes *major*: $(\text{red } G)^{**} p q$ **and** $G \subseteq F$
shows $(\text{red } F)^{**} p q$
<proof>

lemma *red-singleton*: $\text{red } \{f\} p q \longleftrightarrow (\exists t. \text{red-single } p q f t)$
<proof>

lemma *red-single-lookup*:
assumes *red-single* $p q f t$
shows $\text{lookup } q (t \oplus \text{lt } f) = 0$
<proof>

lemma *red-single-higher*:
assumes *red-single* $p q f t$
shows $\text{higher } q (t \oplus \text{lt } f) = \text{higher } p (t \oplus \text{lt } f)$
<proof>

lemma *red-single-ord*:
assumes *red-single* $p q f t$
shows $q \prec_p p$
<proof>

lemma *red-single-nonzero1*:
assumes *red-single* $p q f t$
shows $p \neq 0$
<proof>

lemma *red-single-nonzero2*:
assumes *red-single* $p q f t$
shows $f \neq 0$
<proof>

lemma *red-single-self*:
assumes $p \neq 0$
shows *red-single* $p 0 p 0$
<proof>

lemma *red-single-trans*:
assumes *red-single* $p p0 f t$ **and** $\text{lt } g \text{ adds}_t \text{lt } f$ **and** $g \neq 0$
obtains $p1$ **where** *red-single* $p p1 g (t + (\text{lp } f - \text{lp } g))$
<proof>

lemma *red-nonzero*:
assumes $\text{red } F \ p \ q$
shows $p \neq 0$
 $\langle \text{proof} \rangle$

lemma *red-self*:
assumes $p \neq 0$
shows $\text{red } \{p\} \ p \ 0$
 $\langle \text{proof} \rangle$

lemma *red-ord*:
assumes $\text{red } F \ p \ q$
shows $q \prec_p p$
 $\langle \text{proof} \rangle$

lemma *red-indI1*:
assumes $f \in F$ **and** $f \neq 0$ **and** $p \neq 0$ **and** $\text{adds: } \text{lt } f \ \text{adds}_t \ \text{lt } p$
shows $\text{red } F \ p \ (p - \text{monom-mult } (lc \ p \ / \ lc \ f) \ (lp \ p - lp \ f) \ f)$
 $\langle \text{proof} \rangle$

lemma *red-indI2*:
assumes $p \neq 0$ **and** $r: \text{red } F \ (\text{tail } p) \ q$
shows $\text{red } F \ p \ (q + \text{monomial } (lc \ p) \ (\text{lt } p))$
 $\langle \text{proof} \rangle$

lemma *red-indE*:
assumes $\text{red } F \ p \ q$
shows $(\exists f \in F. f \neq 0 \wedge \text{lt } f \ \text{adds}_t \ \text{lt } p \wedge$
 $(q = p - \text{monom-mult } (lc \ p \ / \ lc \ f) \ (lp \ p - lp \ f) \ f)) \vee$
 $\text{red } F \ (\text{tail } p) \ (q - \text{monomial } (lc \ p) \ (\text{lt } p))$
 $\langle \text{proof} \rangle$

lemma *is-redI*:
assumes $\text{red } F \ a \ b$
shows $\text{is-red } F \ a$
 $\langle \text{proof} \rangle$

lemma *is-redE*:
assumes $\text{is-red } F \ a$
obtains b **where** $\text{red } F \ a \ b$
 $\langle \text{proof} \rangle$

lemma *is-red-alt*:
shows $\text{is-red } F \ a \longleftrightarrow (\exists b. \text{red } F \ a \ b)$
 $\langle \text{proof} \rangle$

lemma *is-red-singletonI*:
assumes $\text{is-red } F \ q$

obtains p where $p \in F$ and $is-red \{p\} q$
<proof>

lemma $is-red-singletonD$:
assumes $is-red \{p\} q$ and $p \in F$
shows $is-red F q$
<proof>

lemma $is-red-singleton-trans$:
assumes $is-red \{f\} p$ and $lt\ g\ adds_t\ lt\ f$ and $g \neq 0$
shows $is-red \{g\} p$
<proof>

lemma $is-red-singleton-not-0$:
assumes $is-red \{f\} p$
shows $f \neq 0$
<proof>

lemma $irred-0$:
shows $\neg is-red F 0$
<proof>

lemma $is-red-indI1$:
assumes $f \in F$ and $f \neq 0$ and $p \neq 0$ and $lt\ f\ adds_t\ lt\ p$
shows $is-red F p$
<proof>

lemma $is-red-indI2$:
assumes $p \neq 0$ and $is-red F (tail\ p)$
shows $is-red F p$
<proof>

lemma $is-red-indE$:
assumes $is-red F p$
shows $(\exists f \in F. f \neq 0 \wedge lt\ f\ adds_t\ lt\ p) \vee is-red F (tail\ p)$
<proof>

lemma $rtrancl-0$:
assumes $(red\ F)\ 0\ x$**
shows $x = 0$
<proof>

lemma $red-rtrancl-ord$:
assumes $(red\ F)\ p\ q$**
shows $q \preceq_p p$
<proof>

lemma $components-red-subset$:
assumes $red F p q$

shows *component-of-term* ' keys $q \subseteq$ *component-of-term* ' keys $p \cup$ *component-of-term* ' Keys F
 ⟨proof⟩

corollary *components-red-rtrancl-subset*:

assumes $(red\ F)^{**}\ p\ q$
shows *component-of-term* ' keys $q \subseteq$ *component-of-term* ' keys $p \cup$ *component-of-term* ' Keys F
 ⟨proof⟩

4.2 Reducibility and Addition & Multiplication

lemma *red-single-monom-mult*:

assumes *red-single* $p\ q\ f\ t$ **and** $c \neq 0$
shows *red-single* $(monom-mult\ c\ s\ p)\ (monom-mult\ c\ s\ q)\ f\ (s + t)$
 ⟨proof⟩

lemma *red-single-plus-1*:

assumes *red-single* $p\ q\ f\ t$ **and** $t \oplus lt\ f \notin keys\ (p + r)$
shows *red-single* $(q + r)\ (p + r)\ f\ t$
 ⟨proof⟩

lemma *red-single-plus-2*:

assumes *red-single* $p\ q\ f\ t$ **and** $t \oplus lt\ f \notin keys\ (q + r)$
shows *red-single* $(p + r)\ (q + r)\ f\ t$
 ⟨proof⟩

lemma *red-single-plus-3*:

assumes *red-single* $p\ q\ f\ t$ **and** $t \oplus lt\ f \in keys\ (p + r)$ **and** $t \oplus lt\ f \in keys\ (q + r)$
shows $\exists s. red-single\ (p + r)\ s\ f\ t \wedge red-single\ (q + r)\ s\ f\ t$
 ⟨proof⟩

lemma *red-single-plus*:

assumes *red-single* $p\ q\ f\ t$
shows *red-single* $(p + r)\ (q + r)\ f\ t \vee$
red-single $(q + r)\ (p + r)\ f\ t \vee$
 $(\exists s. red-single\ (p + r)\ s\ f\ t \wedge red-single\ (q + r)\ s\ f\ t)$ **(is ?A \vee ?B \vee ?C)**
 ⟨proof⟩

lemma *red-single-diff*:

assumes *red-single* $(p - q)\ r\ f\ t$
shows *red-single* $p\ (r + q)\ f\ t \vee red-single\ q\ (p - r)\ f\ t \vee$
 $(\exists p'\ q'. red-single\ p\ p'\ f\ t \wedge red-single\ q\ q'\ f\ t \wedge r = p' - q')$ **(is ?A \vee ?B \vee ?C)**
 ⟨proof⟩

lemma *red-monom-mult*:

assumes $a: red\ F\ p\ q$ **and** $c \neq 0$

shows $\text{red } F (\text{monom-mult } c \ s \ p) (\text{monom-mult } c \ s \ q)$
 ⟨proof⟩

lemma *red-plus-keys-disjoint*:
assumes $\text{red } F \ p \ q$ **and** $\text{keys } p \cap \text{keys } r = \{\}$
shows $\text{red } F (p + r) (q + r)$
 ⟨proof⟩

lemma *red-plus*:
assumes $\text{red } F \ p \ q$
obtains s **where** $(\text{red } F)^{**} (p + r) \ s$ **and** $(\text{red } F)^{**} (q + r) \ s$
 ⟨proof⟩

corollary *red-plus-cs*:
assumes $\text{red } F \ p \ q$
shows $\text{relation.cs } (\text{red } F) (p + r) (q + r)$
 ⟨proof⟩

lemma *red-uminus*:
assumes $\text{red } F \ p \ q$
shows $\text{red } F (-p) (-q)$
 ⟨proof⟩

lemma *red-diff*:
assumes $\text{red } F (p - q) \ r$
obtains $p' \ q'$ **where** $(\text{red } F)^{**} p \ p'$ **and** $(\text{red } F)^{**} q \ q'$ **and** $r = p' - q'$
 ⟨proof⟩

lemma *red-diff-rtrancl'*:
assumes $(\text{red } F)^{**} (p - q) \ r$
obtains $p' \ q'$ **where** $(\text{red } F)^{**} p \ p'$ **and** $(\text{red } F)^{**} q \ q'$ **and** $r = p' - q'$
 ⟨proof⟩

lemma *red-diff-rtrancl*:
assumes $(\text{red } F)^{**} (p - q) \ 0$
obtains s **where** $(\text{red } F)^{**} p \ s$ **and** $(\text{red } F)^{**} q \ s$
 ⟨proof⟩

corollary *red-diff-rtrancl-cs*:
assumes $(\text{red } F)^{**} (p - q) \ 0$
shows $\text{relation.cs } (\text{red } F) \ p \ q$
 ⟨proof⟩

4.3 Confluence of Reducibility

lemma *confluent-distinct-aux*:
assumes $r1: \text{red-single } p \ q1 \ f1 \ t1$ **and** $r2: \text{red-single } p \ q2 \ f2 \ t2$
and $t1 \oplus lt \ f1 \prec_t t2 \oplus lt \ f2$ **and** $f1 \in F$ **and** $f2 \in F$
obtains s **where** $(\text{red } F)^{**} q1 \ s$ **and** $(\text{red } F)^{**} q2 \ s$

<proof>

lemma *confluent-distinct*:

assumes $r1$: *red-single* p $q1$ $f1$ $t1$ **and** $r2$: *red-single* p $q2$ $f2$ $t2$
and ne : $t1 \oplus lt\ f1 \neq t2 \oplus lt\ f2$ **and** $f1 \in F$ **and** $f2 \in F$
obtains s **where** $(red\ F)^{**}$ $q1$ s **and** $(red\ F)^{**}$ $q2$ s

<proof>

corollary *confluent-same*:

assumes $r1$: *red-single* p $q1$ f $t1$ **and** $r2$: *red-single* p $q2$ f $t2$ **and** $f \in F$
obtains s **where** $(red\ F)^{**}$ $q1$ s **and** $(red\ F)^{**}$ $q2$ s

<proof>

4.4 Reducibility and Module Membership

lemma *srtc-in-pmdl*:

assumes *relation.srtc* $(red\ F)$ p q
shows $p - q \in pmdl\ F$

<proof>

lemma *in-pmdl-srtc*:

assumes $p \in pmdl\ F$
shows *relation.srtc* $(red\ F)$ p 0

<proof>

lemma *red-rtranclp-diff-in-pmdl*:

assumes $(red\ F)^{**}$ p q
shows $p - q \in pmdl\ F$

<proof>

corollary *red-diff-in-pmdl*:

assumes $red\ F$ p q
shows $p - q \in pmdl\ F$

<proof>

corollary *red-rtranclp-0-in-pmdl*:

assumes $(red\ F)^{**}$ p 0
shows $p \in pmdl\ F$

<proof>

lemma *pmdl-closed-red*:

assumes $pmdl\ B \subseteq pmdl\ A$ **and** $p \in pmdl\ A$ **and** $red\ B$ p q
shows $q \in pmdl\ A$

<proof>

4.5 More Properties of *red*, *red-single* and *is-red*

lemma *red-rtrancl-mult*:

assumes $(red\ F)^{**}$ p q
shows $(red\ F)^{**}$ $(monom-mult\ c\ t\ p)$ $(monom-mult\ c\ t\ q)$

<proof>

corollary *red-rtrancl-uminus*:

assumes $(red\ F)^{**}\ p\ q$
shows $(red\ F)^{**}\ (-p)\ (-q)$
<proof>

lemma *red-rtrancl-diff-induct* [*consumes 1, case-names base step*]:

assumes $a: (red\ F)^{**}\ (p - q)\ r$
and cases: $P\ p\ p\ !!y\ z. [\ (red\ F)^{**}\ (p - q)\ z; red\ F\ z\ y; P\ p\ (q + z)] ==> P$
 $p\ (q + y)$
shows $P\ p\ (q + r)$
<proof>

lemma *red-rtrancl-diff-0-induct* [*consumes 1, case-names base step*]:

assumes $a: (red\ F)^{**}\ (p - q)\ 0$
and base: $P\ p\ p$ **and** *ind:* $\bigwedge y\ z. [\ (red\ F)^{**}\ (p - q)\ y; red\ F\ y\ z; P\ p\ (y + q)]$
 $==> P\ p\ (z + q)$
shows $P\ p\ q$
<proof>

lemma *is-red-union*: $is-red\ (A\ \cup\ B)\ p\ \longleftrightarrow\ (is-red\ A\ p\ \vee\ is-red\ B\ p)$

<proof>

lemma *red-single-0-lt*:

assumes *red-single* $f\ 0\ h\ t$
shows $lt\ f = t \oplus lt\ h$
<proof>

lemma *red-single-lt-distinct-lt*:

assumes *rs:* *red-single* $f\ g\ h\ t$ **and** $g \neq 0$ **and** $lt\ g \neq lt\ f$
shows $lt\ f = t \oplus lt\ h$
<proof>

lemma *zero-reducibility-implies-lt-divisibility'*:

assumes $(red\ F)^{**}\ f\ 0$ **and** $f \neq 0$
shows $\exists h \in F. h \neq 0 \wedge (lt\ h\ adds_t\ lt\ f)$
<proof>

lemma *zero-reducibility-implies-lt-divisibility*:

assumes $(red\ F)^{**}\ f\ 0$ **and** $f \neq 0$
obtains h **where** $h \in F$ **and** $h \neq 0$ **and** $lt\ h\ adds_t\ lt\ f$
<proof>

lemma *is-red-addsI*:

assumes $f \in F$ **and** $f \neq 0$ **and** $v \in keys\ p$ **and** $lt\ f\ adds_t\ v$
shows $is-red\ F\ p$
<proof>

lemma *is-red-addsE'*:

assumes *is-red F p*

shows $\exists f \in F. \exists v \in \text{keys } p. f \neq 0 \wedge \text{lt } f \text{ adds}_t v$

<proof>

lemma *is-red-addsE*:

assumes *is-red F p*

obtains *f v* **where** $f \in F$ **and** $v \in \text{keys } p$ **and** $f \neq 0$ **and** $\text{lt } f \text{ adds}_t v$

<proof>

lemma *is-red-adds-iff*:

shows $(\text{is-red } F p) \longleftrightarrow (\exists f \in F. \exists v \in \text{keys } p. f \neq 0 \wedge \text{lt } f \text{ adds}_t v)$

<proof>

lemma *is-red-subset*:

assumes *red: is-red A p* **and** *sub: A \subseteq B*

shows *is-red B p*

<proof>

lemma *not-is-red-empty*: $\neg \text{is-red } \{ \} f$

<proof>

lemma *red-single-mult-const*:

assumes *red-single p q f t* **and** $c \neq 0$

shows *red-single p q (monom-mult c 0 f) t*

<proof>

lemma *red-rtrancl-plus-higher*:

assumes $(\text{red } F)^{**} p q$ **and** $\bigwedge u v. u \in \text{keys } p \implies v \in \text{keys } r \implies u \prec_t v$

shows $(\text{red } F)^{**} (p + r) (q + r)$

<proof>

lemma *red-mult-scalar-leading-monomial*: $(\text{red } \{f\})^{**} (p \odot \text{monomial } (lc f) (\text{lt } f))$

$(- p \odot \text{tail } f)$

<proof>

corollary *red-mult-scalar-lt*:

assumes $f \neq 0$

shows $(\text{red } \{f\})^{**} (p \odot \text{monomial } c (\text{lt } f)) (\text{monom-mult } (- c / lc f) 0 (p \odot \text{tail } f))$

<proof>

lemma *is-red-monomial-iff*: $\text{is-red } F (\text{monomial } c v) \longleftrightarrow (c \neq 0 \wedge (\exists f \in F. f \neq 0 \wedge \text{lt } f \text{ adds}_t v))$

<proof>

lemma *is-red-monomialI*:

assumes $c \neq 0$ **and** $f \in F$ **and** $f \neq 0$ **and** $\text{lt } f \text{ adds}_t v$

shows *is-red F (monomial c v)*

$\langle \text{proof} \rangle$

lemma *is-red-monomialD*:

assumes *is-red* F (*monomial* c v)

shows $c \neq 0$

$\langle \text{proof} \rangle$

lemma *is-red-monomialE*:

assumes *is-red* F (*monomial* c v)

obtains f **where** $f \in F$ **and** $f \neq 0$ **and** $\text{lt } f \text{ adds}_t v$

$\langle \text{proof} \rangle$

lemma *replace-lt-adds-stable-is-red*:

assumes *red*: *is-red* F f **and** $q \neq 0$ **and** $\text{lt } q \text{ adds}_t \text{lt } p$

shows *is-red* (*insert* q ($F - \{p\}$)) f

$\langle \text{proof} \rangle$

lemma *conversion-property*:

assumes *is-red* $\{p\}$ f **and** *red* $\{r\}$ p q

shows *is-red* $\{q\}$ $f \vee \text{is-red } \{r\}$ f

$\langle \text{proof} \rangle$

lemma *replace-red-stable-is-red*:

assumes $a1$: *is-red* F f **and** $a2$: *red* ($F - \{p\}$) p q

shows *is-red* (*insert* q ($F - \{p\}$)) f (**is** *is-red* $?F'$ f)

$\langle \text{proof} \rangle$

lemma *is-red-map-scale*:

assumes *is-red* F ($c \cdot p$)

shows *is-red* F p

$\langle \text{proof} \rangle$

corollary *is-irred-map-scale*: $\neg \text{is-red } F$ $p \implies \neg \text{is-red } F$ ($c \cdot p$)

$\langle \text{proof} \rangle$

lemma *is-red-map-scale-iff*: *is-red* F ($c \cdot p$) $\longleftrightarrow (c \neq 0 \wedge \text{is-red } F$ $p)$

$\langle \text{proof} \rangle$

lemma *is-red-uminus*: *is-red* F ($- p$) $\longleftrightarrow \text{is-red } F$ p

$\langle \text{proof} \rangle$

lemma *is-red-plus*:

assumes *is-red* F ($p + q$)

shows *is-red* F $p \vee \text{is-red } F$ q

$\langle \text{proof} \rangle$

lemma *is-irred-plus*: $\neg \text{is-red } F$ $p \implies \neg \text{is-red } F$ $q \implies \neg \text{is-red } F$ ($p + q$)

$\langle \text{proof} \rangle$

lemma *is-red-minus*:
assumes *is-red F (p - q)*
shows *is-red F p \vee is-red F q*
 \langle *proof* \rangle

lemma *is-irred-minus*: \neg *is-red F p* \implies \neg *is-red F q* \implies \neg *is-red F (p - q)*
 \langle *proof* \rangle

end

4.6 Well-foundedness and Termination

context *gd-term*
begin

lemma *dgrad-set-le-red-single*:
assumes *dickson-grading d* **and** *red-single p q f t*
shows *dgrad-set-le d {t}* (*pp-of-term ' keys p*)
 \langle *proof* \rangle

lemma *dgrad-p-set-le-red-single*:
assumes *dickson-grading d* **and** *red-single p q f t*
shows *dgrad-p-set-le d {q} {f, p}*
 \langle *proof* \rangle

lemma *dgrad-p-set-le-red*:
assumes *dickson-grading d* **and** *red F p q*
shows *dgrad-p-set-le d {q}* (*insert p F*)
 \langle *proof* \rangle

corollary *dgrad-p-set-le-red-rtrancl*:
assumes *dickson-grading d* **and** $(red F)^{**}$ *p q*
shows *dgrad-p-set-le d {q}* (*insert p F*)
 \langle *proof* \rangle

lemma *dgrad-p-set-red-single-pp*:
assumes *dickson-grading d* **and** $p \in dgrad-p-set d m$ **and** *red-single p q f t*
shows $d t \leq m$
 \langle *proof* \rangle

lemma *dgrad-p-set-closed-red-single*:
assumes *dickson-grading d* **and** $p \in dgrad-p-set d m$ **and** $f \in dgrad-p-set d m$
and *red-single p q f t*
shows $q \in dgrad-p-set d m$
 \langle *proof* \rangle

lemma *dgrad-p-set-closed-red*:
assumes *dickson-grading d* **and** $F \subseteq dgrad-p-set d m$ **and** $p \in dgrad-p-set d m$
and *red F p q*

shows $q \in \text{dgrad-p-set } d \ m$
 ⟨proof⟩

lemma *dgrad-p-set-closed-red-rtrancl*:

assumes *dickson-grading* d **and** $F \subseteq \text{dgrad-p-set } d \ m$ **and** $p \in \text{dgrad-p-set } d \ m$
and $(\text{red } F)^{**} \ p \ q$
shows $q \in \text{dgrad-p-set } d \ m$
 ⟨proof⟩

lemma *red-rtrancl-repE*:

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$ **and** *finite* G **and** $p \in \text{dgrad-p-set } d \ m$
and $(\text{red } G)^{**} \ p \ r$
obtains q **where** $p = r + (\sum_{g \in G}. q \ g \odot g)$ **and** $\bigwedge g. q \ g \in \text{punit.dgrad-p-set } d \ m$
and $\bigwedge g. \text{lt } (q \ g \odot g) \preceq_t \text{lt } p$
 ⟨proof⟩

lemma *is-relation-order-red*:

assumes *dickson-grading* d
shows *Confluence.relation-order* $(\text{red } F) \ (\prec_p) \ (\text{dgrad-p-set } d \ m)$
 ⟨proof⟩

lemma *red-wf-dgrad-p-set-aux*:

assumes *dickson-grading* d **and** $F \subseteq \text{dgrad-p-set } d \ m$
shows *wfp-on* $(\text{red } F)^{-1-1} \ (\text{dgrad-p-set } d \ m)$
 ⟨proof⟩

lemma *red-wf-dgrad-p-set*:

assumes *dickson-grading* d **and** $F \subseteq \text{dgrad-p-set } d \ m$
shows *wfP* $(\text{red } F)^{-1-1}$
 ⟨proof⟩

lemmas *red-wf-finite = red-wf-dgrad-p-set*[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemma *cbelow-on-monom-mult*:

assumes *dickson-grading* d **and** $F \subseteq \text{dgrad-p-set } d \ m$ **and** $d \ t \leq m$ **and** $c \neq 0$
and *cbelow-on* $(\text{dgrad-p-set } d \ m) \ (\prec_p) \ z \ (\lambda a \ b. \text{red } F \ a \ b \vee \text{red } F \ b \ a) \ p \ q$
shows *cbelow-on* $(\text{dgrad-p-set } d \ m) \ (\prec_p) \ (\text{monom-mult } c \ t \ z) \ (\lambda a \ b. \text{red } F \ a \ b \vee \text{red } F \ b \ a)$
 $(\text{monom-mult } c \ t \ p) \ (\text{monom-mult } c \ t \ q)$
 ⟨proof⟩

lemma *cbelow-on-monom-mult-monomial*:

assumes $c \neq 0$
and *cbelow-on* $(\text{dgrad-p-set } d \ m) \ (\prec_p) \ (\text{monomial } c' \ v) \ (\lambda a \ b. \text{red } F \ a \ b \vee \text{red } F \ b \ a) \ p \ q$
shows *cbelow-on* $(\text{dgrad-p-set } d \ m) \ (\prec_p) \ (\text{monomial } c \ (t \oplus v)) \ (\lambda a \ b. \text{red } F \ a \ b \vee \text{red } F \ b \ a) \ p \ q$

<proof>

lemma *cbelow-on-plus*:

assumes *dickson-grading* d **and** $F \subseteq \text{dgrad-p-set } d \ m$ **and** $r \in \text{dgrad-p-set } d \ m$
and $\text{keys } r \cap \text{keys } z = \{\}$
and $\text{cbelow-on } (\text{dgrad-p-set } d \ m) (\prec_p) z (\lambda a \ b. \text{red } F \ a \ b \vee \text{red } F \ b \ a) \ p \ q$
shows $\text{cbelow-on } (\text{dgrad-p-set } d \ m) (\prec_p) (z + r) (\lambda a \ b. \text{red } F \ a \ b \vee \text{red } F \ b \ a)$
 $(p + r) (q + r)$
<proof>

lemma *is-full-pmdlI-lt-dgrad-p-set*:

assumes *dickson-grading* d **and** $B \subseteq \text{dgrad-p-set } d \ m$
assumes $\bigwedge k. k \in \text{component-of-term } \text{'Keys } (B::('t \Rightarrow_0 'b)::\text{field}) \ \text{set}) \Rightarrow$
 $(\exists b \in B. b \neq 0 \wedge \text{component-of-term } (\text{lt } b) = k \wedge \text{lp } b = 0)$
shows *is-full-pmdl* B
<proof>

lemmas *is-full-pmdlI-lt-finite = is-full-pmdlI-lt-dgrad-p-set*[*OF dickson-grading-dgrad-dummy*
dgrad-p-set-exhaust-expl]

end

4.7 Algorithms

4.7.1 Function *find-adds*

context *ordered-term*
begin

primrec *find-adds* :: $('t \Rightarrow_0 'b) \ \text{list} \Rightarrow 't \Rightarrow ('t \Rightarrow_0 'b::\text{zero}) \ \text{option}$ **where**
 $\text{find-adds } [] \ - = \text{None}$
 $\text{find-adds } (f \ # \ fs) \ u = (\text{if } f \neq 0 \wedge \text{lt } f \ \text{adds}_t \ u \ \text{then } \text{Some } f \ \text{else } \text{find-adds } fs \ u)$

lemma *find-adds-SomeD1*:

assumes $\text{find-adds } fs \ u = \text{Some } f$
shows $f \in \text{set } fs$
<proof>

lemma *find-adds-SomeD2*:

assumes $\text{find-adds } fs \ u = \text{Some } f$
shows $f \neq 0$
<proof>

lemma *find-adds-SomeD3*:

assumes $\text{find-adds } fs \ u = \text{Some } f$
shows $\text{lt } f \ \text{adds}_t \ u$
<proof>

lemma *find-adds-NoneE*:

assumes $\text{find-adds } fs \ u = \text{None}$ **and** $f \in \text{set } fs$

assumes $f = 0 \implies \text{thesis}$ **and** $f \neq 0 \implies \neg \text{lt } f \text{ adds}_t u \implies \text{thesis}$
shows *thesis*
 $\langle \text{proof} \rangle$

lemma *find-adds-SomeD-red-single*:

assumes $p \neq 0$ **and** $\text{find-adds } fs \text{ (lt } p) = \text{Some } f$
shows $\text{red-single } p \text{ (tail } p - \text{monom-mult (lc } p / \text{lc } f) \text{ (lp } p - \text{lp } f) \text{ (tail } f)) } f \text{ (lp } p - \text{lp } f)$
 $\langle \text{proof} \rangle$

lemma *find-adds-SomeD-red*:

assumes $p \neq 0$ **and** $\text{find-adds } fs \text{ (lt } p) = \text{Some } f$
shows $\text{red (set } fs) p \text{ (tail } p - \text{monom-mult (lc } p / \text{lc } f) \text{ (lp } p - \text{lp } f) \text{ (tail } f))}$
 $\langle \text{proof} \rangle$

end

4.7.2 Function *trd*

context *gd-term*

begin

definition *trd-term* :: $('a \Rightarrow \text{nat}) \Rightarrow ((('t \Rightarrow_0 'b::\text{field}) \text{list} \times ('t \Rightarrow_0 'b) \times ('t \Rightarrow_0 'b)) \times$

$((('t \Rightarrow_0 'b) \text{list} \times ('t \Rightarrow_0 'b) \times ('t \Rightarrow_0 'b))) \text{set}$

where $\text{trd-term } d = \{(x, y). \text{dgrad-p-set-le } d \text{ (set (fst (snd } x) \# \text{fst } x)) \text{ (set (fst (snd } y) \# \text{fst } y))} \wedge \text{fst (snd } x) \prec_p \text{fst (snd } y)\}$

lemma *trd-term-wf*:

assumes *dickson-grading* d
shows $\text{wf (trd-term } d)$
 $\langle \text{proof} \rangle$

function *trd-aux* :: $('t \Rightarrow_0 'b) \text{list} \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow ('t \Rightarrow_0 'b::\text{field})$

where

$\text{trd-aux } fs \text{ } p \text{ } r =$
 $(\text{if } p = 0 \text{ then}$
 $\quad r$
 else
 $\quad \text{case find-adds } fs \text{ (lt } p) \text{ of}$
 $\quad \quad \text{None} \Rightarrow \text{trd-aux } fs \text{ (tail } p) \text{ (r + monomial (lc } p) \text{ (lt } p))}$
 $\quad \quad | \text{Some } f \Rightarrow \text{trd-aux } fs \text{ (tail } p - \text{monom-mult (lc } p / \text{lc } f) \text{ (lp } p - \text{lp } f) \text{ (tail } f)) } r$
 $\quad)$

$\langle \text{proof} \rangle$

termination $\langle \text{proof} \rangle$

definition *trd* :: $('t \Rightarrow_0 'b::\text{field}) \text{list} \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow ('t \Rightarrow_0 'b)$

where $\text{trd } fs \text{ } p = \text{trd-aux } fs \text{ } p \text{ } 0$

lemma *trd-aux-red-rtrancl*: $(\text{red } (\text{set } fs))^{**} p (\text{trd-aux } fs p r - r)$
 <proof>

corollary *trd-red-rtrancl*: $(\text{red } (\text{set } fs))^{**} p (\text{trd } fs p)$
 <proof>

lemma *trd-aux-irred*:
assumes $\neg \text{is-red } (\text{set } fs) r$
shows $\neg \text{is-red } (\text{set } fs) (\text{trd-aux } fs p r)$
 <proof>

corollary *trd-irred*: $\neg \text{is-red } (\text{set } fs) (\text{trd } fs p)$
 <proof>

lemma *trd-in-pmdl*: $p - (\text{trd } fs p) \in \text{pmdl } (\text{set } fs)$
 <proof>

lemma *pmdl-closed-trd*:
assumes $p \in \text{pmdl } B$ **and** $\text{set } fs \subseteq \text{pmdl } B$
shows $(\text{trd } fs p) \in \text{pmdl } B$
 <proof>

end

end

5 Gröbner Bases and Buchberger's Theorem

theory *Groebner-Bases*
imports *Reduction*
begin

This theory provides the main results about Gröbner bases for modules of multivariate polynomials.

context *gd-term*
begin

definition *crit-pair* :: $(t \Rightarrow_0 'b::\text{field}) \Rightarrow (t \Rightarrow_0 'b) \Rightarrow ((t \Rightarrow_0 'b) \times (t \Rightarrow_0 'b))$
where *crit-pair* $p q =$
 (if *component-of-term* $(lt p) = \text{component-of-term } (lt q)$ then
 (*monom-mult* $(1 / lc p) ((lcs (lp p) (lp q)) - (lp p)) (\text{tail } p)$,
monom-mult $(1 / lc q) ((lcs (lp p) (lp q)) - (lp q)) (\text{tail } q)$)
 else $(0, 0)$)

definition *crit-pair-cbelow-on* :: $(a \Rightarrow \text{nat}) \Rightarrow \text{nat} \Rightarrow (t \Rightarrow_0 'b::\text{field}) \text{ set} \Rightarrow (t \Rightarrow_0 'b) \Rightarrow (t \Rightarrow_0 'b) \Rightarrow \text{bool}$
where *crit-pair-cbelow-on* $d m F p q \longleftrightarrow$
cbelow-on $(dgrad-p\text{-set } d m) (\prec_p)$

(*monomial 1 (term-of-pair (lcs (lp p) (lp q), component-of-term (lt p))))*
(λa b. red F a b ∨ red F b a) (fst (crit-pair p q)) (snd (crit-pair p q))

definition *spoly* :: (*t* ⇒₀ 'b) ⇒ (*t* ⇒₀ 'b) ⇒ (*t* ⇒₀ 'b::field)
where *spoly p q* = (*let v1 = lt p; v2 = lt q in*
if component-of-term v1 = component-of-term v2 then
let t1 = pp-of-term v1; t2 = pp-of-term v2; l = lcs t1 t2 in
(monom-mult (1 / lookup p v1) (l - t1) p) - (monom-mult (1 / lookup q v2) (l - t2) q)
else 0)

definition (*in ordered-term*) *is-Groebner-basis* :: (*t* ⇒₀ 'b::field) set ⇒ bool
where *is-Groebner-basis F* ≡ *relation.is-ChurchRosser (red F)*

5.1 Critical Pairs and S-Polynomials

lemma *crit-pair-same*: *fst (crit-pair p p) = snd (crit-pair p p)*
<proof>

lemma *crit-pair-swap*: *crit-pair p q = (snd (crit-pair q p), fst (crit-pair q p))*
<proof>

lemma *crit-pair-zero [simp]*: *fst (crit-pair 0 q) = 0 and snd (crit-pair p 0) = 0*
<proof>

lemma *dgrad-p-set-le-crit-pair-zero*: *dgrad-p-set-le d {fst (crit-pair p 0)} {p}*
<proof>

lemma *dgrad-p-set-le-fst-crit-pair*:
assumes *dickson-grading d*
shows *dgrad-p-set-le d {fst (crit-pair p q)} {p, q}*
<proof>

lemma *dgrad-p-set-le-snd-crit-pair*:
assumes *dickson-grading d*
shows *dgrad-p-set-le d {snd (crit-pair p q)} {p, q}*
<proof>

lemma *dgrad-p-set-closed-fst-crit-pair*:
assumes *dickson-grading d and p ∈ dgrad-p-set d m and q ∈ dgrad-p-set d m*
shows *fst (crit-pair p q) ∈ dgrad-p-set d m*
<proof>

lemma *dgrad-p-set-closed-snd-crit-pair*:
assumes *dickson-grading d and p ∈ dgrad-p-set d m and q ∈ dgrad-p-set d m*
shows *snd (crit-pair p q) ∈ dgrad-p-set d m*
<proof>

lemma *fst-crit-pair-below-lcs:*

fst (crit-pair p q) \prec_p monomial 1 (term-of-pair (lcs (lp p) (lp q), component-of-term (lt p)))
<proof>

lemma *snd-crit-pair-below-lcs:*

snd (crit-pair p q) \prec_p monomial 1 (term-of-pair (lcs (lp p) (lp q), component-of-term (lt p)))
<proof>

lemma *crit-pair-cbelow-same:*

assumes dickson-grading d and $p \in \text{dgrad-p-set } d \ m$
shows crit-pair-cbelow-on d m F p p
<proof>

lemma *crit-pair-cbelow-distinct-component:*

assumes component-of-term (lt p) \neq component-of-term (lt q)
shows crit-pair-cbelow-on d m F p q
<proof>

lemma *crit-pair-cbelow-sym:*

assumes crit-pair-cbelow-on d m F p q
shows crit-pair-cbelow-on d m F q p
<proof>

lemma *crit-pair-cs-imp-crit-pair-cbelow-on:*

assumes dickson-grading d and $F \subseteq \text{dgrad-p-set } d \ m$ and $p \in \text{dgrad-p-set } d \ m$
and $q \in \text{dgrad-p-set } d \ m$
and relation.cs (red F) (fst (crit-pair p q)) (snd (crit-pair p q))
shows crit-pair-cbelow-on d m F p q
<proof>

lemma *crit-pair-cbelow-mono:*

assumes crit-pair-cbelow-on d m F p q and $F \subseteq G$
shows crit-pair-cbelow-on d m G p q
<proof>

lemma *lcs-red-single-fst-crit-pair:*

assumes $p \neq 0$ and component-of-term (lt p) = component-of-term (lt q)
defines t1 \equiv lp p
defines t2 \equiv lp q
shows red-single (monomial (- 1) (term-of-pair (lcs t1 t2, component-of-term (lt p))))
(fst (crit-pair p q)) p (lcs t1 t2 - t1)
<proof>

corollary *lcs-red-single-snd-crit-pair:*

assumes $q \neq 0$ and component-of-term (lt p) = component-of-term (lt q)

defines $t1 \equiv lp\ p$
defines $t2 \equiv lp\ q$
shows $red\text{-}single\ (monomial\ (-\ 1)\ (term\text{-}of\text{-}pair\ (lcs\ t1\ t2,\ component\text{-}of\text{-}term\ (lt\ p))))$
 $(snd\ (crit\text{-}pair\ p\ q))\ q\ (lcs\ t1\ t2\ -\ t2)$
 $\langle proof \rangle$

lemma $GB\text{-}imp\text{-}crit\text{-}pair\text{-}cbelow\text{-}dgrad\text{-}p\text{-}set$:
assumes $dickson\text{-}grading\ d$ **and** $F \subseteq dgrad\text{-}p\text{-}set\ d\ m$ **and** $is\text{-}Groebner\text{-}basis\ F$
assumes $p \in F$ **and** $q \in F$ **and** $p \neq 0$ **and** $q \neq 0$
shows $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ F\ p\ q$
 $\langle proof \rangle$

lemma $spoly\text{-}alt$:
assumes $p \neq 0$ **and** $q \neq 0$
shows $spoly\ p\ q = fst\ (crit\text{-}pair\ p\ q) - snd\ (crit\text{-}pair\ p\ q)$
 $\langle proof \rangle$

lemma $spoly\text{-}same$: $spoly\ p\ p = 0$
 $\langle proof \rangle$

lemma $spoly\text{-}swap$: $spoly\ p\ q = -\ spoly\ q\ p$
 $\langle proof \rangle$

lemma $spoly\text{-}red\text{-}zero\text{-}imp\text{-}crit\text{-}pair\text{-}cbelow\text{-}on$:
assumes $dickson\text{-}grading\ d$ **and** $F \subseteq dgrad\text{-}p\text{-}set\ d\ m$ **and** $p \in dgrad\text{-}p\text{-}set\ d\ m$
and $q \in dgrad\text{-}p\text{-}set\ d\ m$ **and** $p \neq 0$ **and** $q \neq 0$ **and** $(red\ F)^{**}\ (spoly\ p\ q)\ 0$
shows $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ F\ p\ q$
 $\langle proof \rangle$

lemma $dgrad\text{-}p\text{-}set\text{-}le\text{-}spoly\text{-}zero$: $dgrad\text{-}p\text{-}set\text{-}le\ d\ \{spoly\ p\ 0\}\ \{p\}$
 $\langle proof \rangle$

lemma $dgrad\text{-}p\text{-}set\text{-}le\text{-}spoly$:
assumes $dickson\text{-}grading\ d$
shows $dgrad\text{-}p\text{-}set\text{-}le\ d\ \{spoly\ p\ q\}\ \{p,\ q\}$
 $\langle proof \rangle$

lemma $dgrad\text{-}p\text{-}set\text{-}closed\text{-}spoly$:
assumes $dickson\text{-}grading\ d$ **and** $p \in dgrad\text{-}p\text{-}set\ d\ m$ **and** $q \in dgrad\text{-}p\text{-}set\ d\ m$
shows $spoly\ p\ q \in dgrad\text{-}p\text{-}set\ d\ m$
 $\langle proof \rangle$

lemma $components\text{-}spoly\text{-}subset$: $component\text{-}of\text{-}term\ 'keys\ (spoly\ p\ q) \subseteq component\text{-}of\text{-}term\ 'Keys\ \{p,\ q\}$
 $\langle proof \rangle$

lemma $pmdl\text{-}closed\text{-}spoly$:
assumes $p \in pmdl\ F$ **and** $q \in pmdl\ F$

shows $\text{spoly } p \ q \in \text{pmdl } F$
 ⟨proof⟩

5.2 Buchberger's Theorem

Before proving the main theorem of Gröbner bases theory for S-polynomials, as is usually done in textbooks, we first prove it for critical pairs: a set F yields a confluent reduction relation if the critical pairs of all $p \in F$ and $q \in F$ can be connected below the least common sum of the leading power-products of p and q . The reason why we proceed in this way is that it becomes much easier to prove the correctness of Buchberger's second criterion for avoiding useless pairs.

lemma *crit-pair-cbelow-imp-confluent-dgrad-p-set:*

assumes $\text{dg: dickson-grading } d$ **and** $F \subseteq \text{dgrad-p-set } d \ m$

assumes main: $\bigwedge p \ q. p \in F \implies q \in F \implies p \neq 0 \implies q \neq 0 \implies \text{crit-pair-cbelow-on } d \ m \ F \ p \ q$

shows $\text{relation.is-confluent-on } (\text{red } F) \ (\text{dgrad-p-set } d \ m)$

⟨proof⟩

corollary *crit-pair-cbelow-imp-GB-dgrad-p-set:*

assumes $\text{dickson-grading } d$ **and** $F \subseteq \text{dgrad-p-set } d \ m$

assumes $\bigwedge p \ q. p \in F \implies q \in F \implies p \neq 0 \implies q \neq 0 \implies \text{crit-pair-cbelow-on } d \ m \ F \ p \ q$

shows $\text{is-Groebner-basis } F$

⟨proof⟩

corollary *Buchberger-criterion-dgrad-p-set:*

assumes $\text{dickson-grading } d$ **and** $F \subseteq \text{dgrad-p-set } d \ m$

assumes $\bigwedge p \ q. p \in F \implies q \in F \implies p \neq 0 \implies q \neq 0 \implies p \neq q \implies \text{component-of-term } (\text{lt } p) = \text{component-of-term } (\text{lt } q) \implies (\text{red } F)** \ (\text{spoly } p \ q) \ 0$

shows $\text{is-Groebner-basis } F$

⟨proof⟩

lemmas $\text{Buchberger-criterion-finite} = \text{Buchberger-criterion-dgrad-p-set}[\text{OF } \text{dickson-grading-dgrad-dummy } \text{dgrad-p-set-exhaust-expl}]$

lemma (in *ordered-term*) *GB-imp-zero-reducibility:*

assumes $\text{is-Groebner-basis } G$ **and** $f \in \text{pmdl } G$

shows $(\text{red } G)** \ f \ 0$

⟨proof⟩

lemma (in *ordered-term*) *GB-imp-reducibility:*

assumes $\text{is-Groebner-basis } G$ **and** $f \neq 0$ **and** $f \in \text{pmdl } G$

shows $\text{is-red } G \ f$

⟨proof⟩

lemma *is-Groebner-basis-empty:* $\text{is-Groebner-basis } \{\}$

<proof>

lemma *is-Groebner-basis-singleton: is-Groebner-basis {f}*
<proof>

5.3 Buchberger's Criteria for Avoiding Useless Pairs

Unfortunately, the product criterion is only applicable to scalar polynomials.

lemma (in *gd-powerprod*) *product-criterion:*

assumes *dickson-grading d and $F \subseteq \text{punit.dgrad-p-set } d \ m$ and $p \in F$ and $q \in F$*

and *$p \neq 0$ and $q \neq 0$ and $\text{gcs } (\text{punit.lt } p) (\text{punit.lt } q) = 0$*

shows *$\text{punit.crit-pair-cbelow-on } d \ m \ F \ p \ q$*
<proof>

lemma *chain-criterion:*

assumes *dickson-grading d and $F \subseteq \text{dgrad-p-set } d \ m$ and $p \in F$ and $q \in F$*

and *$p \neq 0$ and $q \neq 0$ and $\text{lp } r \text{ adds lcs } (\text{lp } p) (\text{lp } q)$*

and *$\text{component-of-term } (\text{lt } r) = \text{component-of-term } (\text{lt } p)$*

and *$\text{pr: crit-pair-cbelow-on } d \ m \ F \ p \ r$ and $\text{rq: crit-pair-cbelow-on } d \ m \ F \ r \ q$*

shows *$\text{crit-pair-cbelow-on } d \ m \ F \ p \ q$*
<proof>

5.4 Weak and Strong Gröbner Bases

lemma *ord-p-wf-on:*

assumes *dickson-grading d*

shows *$\text{wfp-on } (\prec_p) (\text{dgrad-p-set } d \ m)$*

<proof>

lemma *is-red-implies-0-red-dgrad-p-set:*

assumes *dickson-grading d and $B \subseteq \text{dgrad-p-set } d \ m$*

assumes *$\text{pmdl } B \subseteq \text{pmdl } A$ and $\bigwedge q. q \in \text{pmdl } A \implies q \in \text{dgrad-p-set } d \ m \implies q \neq 0 \implies \text{is-red } B \ q$*

and *$p \in \text{pmdl } A$ and $p \in \text{dgrad-p-set } d \ m$*

shows *$(\text{red } B)^{**} \ p \ 0$*
<proof>

lemma *is-red-implies-0-red-dgrad-p-set':*

assumes *dickson-grading d and $B \subseteq \text{dgrad-p-set } d \ m$*

assumes *$\text{pmdl } B \subseteq \text{pmdl } A$ and $\bigwedge q. q \in \text{pmdl } A \implies q \neq 0 \implies \text{is-red } B \ q$*

and *$p \in \text{pmdl } A$*

shows *$(\text{red } B)^{**} \ p \ 0$*
<proof>

lemma *pmdl-eqI-adds-lt-dgrad-p-set:*

fixes *$G::('t \Rightarrow_0 'b::\text{field}) \text{ set}$*

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$ **and** $B \subseteq \text{dgrad-p-set } d \ m$
and $\text{pmdl } G \subseteq \text{pmdl } B$
assumes $\bigwedge f. f \in \text{pmdl } B \implies f \in \text{dgrad-p-set } d \ m \implies f \neq 0 \implies (\exists g \in G. g \neq 0 \wedge \text{lt } g \text{ adds}_t \text{ lt } f)$
shows $\text{pmdl } G = \text{pmdl } B$
 $\langle \text{proof} \rangle$

lemma *pmdl-eqI-adds-lt-dgrad-p-set'*:
fixes $G::('t \Rightarrow_0 'b::\text{field}) \text{ set}$
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$ **and** $\text{pmdl } G \subseteq \text{pmdl } B$
assumes $\bigwedge f. f \in \text{pmdl } B \implies f \neq 0 \implies (\exists g \in G. g \neq 0 \wedge \text{lt } g \text{ adds}_t \text{ lt } f)$
shows $\text{pmdl } G = \text{pmdl } B$
 $\langle \text{proof} \rangle$

lemma *GB-implies-unique-nf-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes *isGB: is-Groebner-basis* G
shows $\exists! h. (\text{red } G)^{**} f \ h \wedge \neg \text{is-red } G \ h$
 $\langle \text{proof} \rangle$

lemma *translation-property'*:
assumes $p \neq 0$ **and** *red-p-0*: $(\text{red } F)^{**} p \ 0$
shows $\text{is-red } F \ (p + q) \vee \text{is-red } F \ q$
 $\langle \text{proof} \rangle$

lemma *translation-property*:
assumes $p \neq q$ **and** *red-0*: $(\text{red } F)^{**} (p - q) \ 0$
shows $\text{is-red } F \ p \vee \text{is-red } F \ q$
 $\langle \text{proof} \rangle$

lemma *weak-GB-is-strong-GB-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes $\bigwedge f. f \in \text{pmdl } G \implies f \in \text{dgrad-p-set } d \ m \implies (\text{red } G)^{**} f \ 0$
shows *is-Groebner-basis* G
 $\langle \text{proof} \rangle$

lemma *weak-GB-is-strong-GB*:
assumes $\bigwedge f. f \in (\text{pmdl } G) \implies (\text{red } G)^{**} f \ 0$
shows *is-Groebner-basis* G
 $\langle \text{proof} \rangle$

corollary *GB-alt-1-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
shows *is-Groebner-basis* $G \iff (\forall f \in \text{pmdl } G. f \in \text{dgrad-p-set } d \ m \implies (\text{red } G)^{**} f \ 0)$
 $\langle \text{proof} \rangle$

corollary *GB-alt-1: is-Groebner-basis* $G \iff (\forall f \in \text{pmdl } G. (\text{red } G)^{**} f \ 0)$
 $\langle \text{proof} \rangle$

lemma *isGB-I-is-red:*

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$

assumes $\bigwedge f. f \in \text{pmdl } G \implies f \in \text{dgrad-p-set } d \ m \implies f \neq 0 \implies \text{is-red } G \ f$

shows *is-Groebner-basis* G

<proof>

lemma *GB-alt-2-dgrad-p-set:*

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$

shows *is-Groebner-basis* $G \longleftrightarrow (\forall f \in \text{pmdl } G. f \neq 0 \longrightarrow \text{is-red } G \ f)$

<proof>

lemma *GB-adds-lt:*

assumes *is-Groebner-basis* G **and** $f \in \text{pmdl } G$ **and** $f \neq 0$

obtains g **where** $g \in G$ **and** $g \neq 0$ **and** $\text{lt } g \ \text{adds}_t \ \text{lt } f$

<proof>

lemma *isGB-I-adds-lt:*

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$

assumes $\bigwedge f. f \in \text{pmdl } G \implies f \in \text{dgrad-p-set } d \ m \implies f \neq 0 \implies (\exists g \in G. g \neq 0 \wedge \text{lt } g \ \text{adds}_t \ \text{lt } f)$

shows *is-Groebner-basis* G

<proof>

lemma *GB-alt-3-dgrad-p-set:*

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$

shows *is-Groebner-basis* $G \longleftrightarrow (\forall f \in \text{pmdl } G. f \neq 0 \longrightarrow (\exists g \in G. g \neq 0 \wedge \text{lt } g \ \text{adds}_t \ \text{lt } f))$

(**is** $?L \longleftrightarrow ?R$)

<proof>

lemma *GB-insert:*

assumes *is-Groebner-basis* G **and** $f \in \text{pmdl } G$

shows *is-Groebner-basis* (*insert* $f \ G$)

<proof>

lemma *GB-subset:*

assumes *is-Groebner-basis* G **and** $G \subseteq G'$ **and** $\text{pmdl } G' = \text{pmdl } G$

shows *is-Groebner-basis* G'

<proof>

lemma (**in** *ordered-term*) *GB-remove-0-stable-GB:*

assumes *is-Groebner-basis* G

shows *is-Groebner-basis* ($G - \{0\}$)

<proof>

lemmas *is-red-implies-0-red-finite = is-red-implies-0-red-dgrad-p-set*[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemmas *GB-implies-unique-nf-finite = GB-implies-unique-nf-dgrad-p-set*[*OF dick-*

son-grading-dgrad-dummy dgrad-p-set-exhaust-expl
lemmas *GB-alt-2-finite = GB-alt-2-dgrad-p-set*[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]
lemmas *GB-alt-3-finite = GB-alt-3-dgrad-p-set*[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]
lemmas *pmdl-eqI-adds-lt-finite = pmdl-eqI-adds-lt-dgrad-p-set*'[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

5.5 Alternative Characterization of Gröbner Bases via Representations of S-Polynomials

definition *spoly-rep* :: (*'a* ⇒ *nat*) ⇒ *nat* ⇒ (*'t* ⇒₀ *'b*) *set* ⇒ (*'t* ⇒₀ *'b*) ⇒ (*'t* ⇒₀ *'b*::*field*) ⇒ *bool*

where *spoly-rep d m G g1 g2* ⇔ (∃ *q*. *spoly g1 g2* = (∑ *g* ∈ *G*. *q g* ⊙ *g*) ∧
(∀ *g*. *q g* ∈ *punit.dgrad-p-set d m* ∧
(*q g* ⊙ *g* ≠ 0 ⇒ *lt (q g* ⊙ *g)* <_{*t*} *term-of-pair (lcs (lp g1) (lp*
g2),
component-of-term (lt g2)))))

lemma *spoly-repI*:

spoly g1 g2 = (∑ *g* ∈ *G*. *q g* ⊙ *g*) ⇒ (∧ *g*. *q g* ∈ *punit.dgrad-p-set d m*) ⇒
(∧ *g*. *q g* ⊙ *g* ≠ 0 ⇒ *lt (q g* ⊙ *g)* <_{*t*} *term-of-pair (lcs (lp g1) (lp g2)*,
component-of-term (lt g2))) ⇒
spoly-rep d m G g1 g2
⟨*proof*⟩

lemma *spoly-repI-zero*:

assumes *spoly g1 g2 = 0*
shows *spoly-rep d m G g1 g2*
⟨*proof*⟩

lemma *spoly-repE*:

assumes *spoly-rep d m G g1 g2*
obtains *q* **where** *spoly g1 g2* = (∑ *g* ∈ *G*. *q g* ⊙ *g*) **and** ∧ *g*. *q g* ∈ *punit.dgrad-p-set d m*
and ∧ *g*. *q g* ⊙ *g* ≠ 0 ⇒ *lt (q g* ⊙ *g)* <_{*t*} *term-of-pair (lcs (lp g1) (lp g2)*,
component-of-term (lt g2))
⟨*proof*⟩

corollary *isGB-D-spoly-rep*:

assumes *dickson-grading d* **and** *is-Groebner-basis G* **and** *G* ⊆ *dgrad-p-set d m*
and *finite G*
and *g1* ∈ *G* **and** *g2* ∈ *G* **and** *g1* ≠ 0 **and** *g2* ≠ 0
shows *spoly-rep d m G g1 g2*
⟨*proof*⟩

The finiteness assumption on *G* in the following theorem could be dropped, but it makes the proof a lot easier (although it is still fairly complicated).

lemma *isGB-I-spoly-rep*:

assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$ **and** *finite* G
and $\bigwedge g1 \ g2. \ g1 \in G \implies g2 \in G \implies g1 \neq 0 \implies g2 \neq 0 \implies \text{spoly } g1 \ g2 \neq 0 \implies \text{spoly-rep } d \ m \ G \ g1 \ g2$
shows *is-Groebner-basis* G
 $\langle \text{proof} \rangle$

5.6 Replacing Elements in Gröbner Bases

lemma *replace-in-dgrad-p-set*:
assumes $G \subseteq \text{dgrad-p-set } d \ m$
obtains n **where** $q \in \text{dgrad-p-set } d \ n$ **and** $G \subseteq \text{dgrad-p-set } d \ n$
and $\text{insert } q \ (G - \{p\}) \subseteq \text{dgrad-p-set } d \ n$
 $\langle \text{proof} \rangle$

lemma *GB-replace-lt-adds-stable-GB-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes *isGB: is-Groebner-basis* G **and** $q \neq 0$ **and** $q: q \in (\text{pmdl } G)$ **and** $lt \ q$
 $\text{adds}_t \ lt \ p$
shows *is-Groebner-basis* $(\text{insert } q \ (G - \{p\}))$ (**is** *is-Groebner-basis* $?G'$)
 $\langle \text{proof} \rangle$

lemma *GB-replace-lt-adds-stable-pmdl-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes *isGB: is-Groebner-basis* G **and** $q \neq 0$ **and** $q \in \text{pmdl } G$ **and** $lt \ q \ \text{adds}_t$
 $lt \ p$
shows $\text{pmdl } (\text{insert } q \ (G - \{p\})) = \text{pmdl } G$ (**is** $\text{pmdl } ?G' = \text{pmdl } G$)
 $\langle \text{proof} \rangle$

lemma *GB-replace-red-stable-GB-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes *isGB: is-Groebner-basis* G **and** $p \in G$ **and** $q: \text{red } (G - \{p\}) \ p \ q$
shows *is-Groebner-basis* $(\text{insert } q \ (G - \{p\}))$ (**is** *is-Groebner-basis* $?G'$)
 $\langle \text{proof} \rangle$

lemma *GB-replace-red-stable-pmdl-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes *isGB: is-Groebner-basis* G **and** $p \in G$ **and** $ptoq: \text{red } (G - \{p\}) \ p \ q$
shows $\text{pmdl } (\text{insert } q \ (G - \{p\})) = \text{pmdl } G$ (**is** $\text{pmdl } ?G' = -$)
 $\langle \text{proof} \rangle$

lemma *GB-replace-red-rtranclp-stable-GB-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$
assumes *isGB: is-Groebner-basis* G **and** $p \in G$ **and** $ptoq: (\text{red } (G - \{p\}))^{**} \ p \ q$
shows *is-Groebner-basis* $(\text{insert } q \ (G - \{p\}))$
 $\langle \text{proof} \rangle$

lemma *GB-replace-red-rtranclp-stable-pmdl-dgrad-p-set*:
assumes *dickson-grading* d **and** $G \subseteq \text{dgrad-p-set } d \ m$

assumes *isGB*: *is-Groebner-basis* G **and** $p \in G$ **and** *ptoq*: $(\text{red } (G - \{p\}))^{**} p$
 q
shows $\text{pmdl } (\text{insert } q (G - \{p\})) = \text{pmdl } G$
 $\langle \text{proof} \rangle$

lemmas *GB-replace-lt-adds-stable-GB-finite* =

GB-replace-lt-adds-stable-GB-dgrad-p-set[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemmas *GB-replace-lt-adds-stable-pmdl-finite* =

GB-replace-lt-adds-stable-pmdl-dgrad-p-set[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemmas *GB-replace-red-stable-GB-finite* =

GB-replace-red-stable-GB-dgrad-p-set[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemmas *GB-replace-red-stable-pmdl-finite* =

GB-replace-red-stable-pmdl-dgrad-p-set[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemmas *GB-replace-red-rtranclp-stable-GB-finite* =

GB-replace-red-rtranclp-stable-GB-dgrad-p-set[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

lemmas *GB-replace-red-rtranclp-stable-pmdl-finite* =

GB-replace-red-rtranclp-stable-pmdl-dgrad-p-set[*OF dickson-grading-dgrad-dummy dgrad-p-set-exhaust-expl*]

5.7 An Inconstructive Proof of the Existence of Finite Gröbner Bases

lemma *ex-finite-GB-dgrad-p-set*:

assumes *dickson-grading* d **and** *finite* (*component-of-term* ‘*Keys* F) **and** $F \subseteq \text{dgrad-p-set } d \ m$

obtains G **where** $G \subseteq \text{dgrad-p-set } d \ m$ **and** *finite* G **and** *is-Groebner-basis* G **and** $\text{pmdl } G = \text{pmdl } F$

$\langle \text{proof} \rangle$

The preceding lemma justifies the following definition.

definition *some-GB* :: $(t \Rightarrow_0 b)$ *set* $\Rightarrow (t \Rightarrow_0 b::\text{field})$ *set*

where *some-GB* $F = (\text{SOME } G. \text{finite } G \wedge \text{is-Groebner-basis } G \wedge \text{pmdl } G = \text{pmdl } F)$

lemma *some-GB-props-dgrad-p-set*:

assumes *dickson-grading* d **and** *finite* (*component-of-term* ‘*Keys* F) **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows $\text{finite } (\text{some-GB } F) \wedge \text{is-Groebner-basis } (\text{some-GB } F) \wedge \text{pmdl } (\text{some-GB } F) = \text{pmdl } F$

$\langle \text{proof} \rangle$

lemma *finite-some-GB-dgrad-p-set*:

assumes *dickson-grading* d **and** *finite* (*component-of-term* ‘*Keys* F) **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows $\text{finite } (\text{some-GB } F)$

$\langle \text{proof} \rangle$

lemma *some-GB-isGB-dgrad-p-set*:

assumes *dickson-grading* d **and** *finite* (*component-of-term* ‘*Keys* F) **and** $F \subseteq$
dgrad-p-set d m
shows *is-Groebner-basis* (*some-GB* F)
 \langle *proof* \rangle

lemma *some-GB-pmdl-dgrad-p-set*:

assumes *dickson-grading* d **and** *finite* (*component-of-term* ‘*Keys* F) **and** $F \subseteq$
dgrad-p-set d m
shows *pmdl* (*some-GB* F) = *pmdl* F
 \langle *proof* \rangle

lemma *finite-imp-finite-component-Keys*:

assumes *finite* F
shows *finite* (*component-of-term* ‘*Keys* F)
 \langle *proof* \rangle

lemma *finite-some-GB-finite*: *finite* $F \implies$ *finite* (*some-GB* F)

\langle *proof* \rangle

lemma *some-GB-isGB-finite*: *finite* $F \implies$ *is-Groebner-basis* (*some-GB* F)

\langle *proof* \rangle

lemma *some-GB-pmdl-finite*: *finite* $F \implies$ *pmdl* (*some-GB* F) = *pmdl* F

\langle *proof* \rangle

Theory *Buchberger* implements an algorithm for effectively computing Gröbner bases.

5.8 Relation *red-supset*

The following relation is needed for proving the termination of Buchberger’s algorithm (i. e. function *gb-schema-aux*).

definition *red-supset*::($'t \Rightarrow_0 'b::field$) *set* \Rightarrow ($'t \Rightarrow_0 'b$) *set* \Rightarrow *bool* (**infixl** $\sqsupset p$ 50)

where *red-supset* A $B \equiv$ ($\exists p. is-red$ A $p \wedge \neg is-red$ B p) \wedge ($\forall p. is-red$ B $p \longrightarrow is-red$ A p)

lemma *red-supsetE*:

assumes $A \sqsupset p$ B
obtains p **where** *is-red* A p **and** $\neg is-red$ B p
 \langle *proof* \rangle

lemma *red-supsetD*:

assumes $a1: A \sqsupset p$ B **and** $a2: is-red$ B p
shows *is-red* A p
 \langle *proof* \rangle

lemma *red-supsetI* [*intro*]:

assumes $\bigwedge q. is-red$ B $q \implies is-red$ A q **and** *is-red* A p **and** $\neg is-red$ B p

shows $A \sqsupset_p B$
 $\langle proof \rangle$

lemma *red-supset-insertI*:
assumes $x \neq 0$ **and** $\neg is-red\ A\ x$
shows $(insert\ x\ A) \sqsupset_p A$
 $\langle proof \rangle$

lemma *red-supset-transitive*:
assumes $A \sqsupset_p B$ **and** $B \sqsupset_p C$
shows $A \sqsupset_p C$
 $\langle proof \rangle$

lemma *red-supset-wf-on*:
assumes *dickson-grading* d **and** *finite* K
shows $wfp-on\ (\sqsupset_p)\ (Pow\ (dgrad-p-set\ d\ m) \cap \{F.\ component-of-term\ 'Keys\ F\} \subseteq K)$
 $\langle proof \rangle$

end

lemma *in-lex-prod-alt*:
 $(x, y) \in r < *lex* > s \longleftrightarrow (((fst\ x), (fst\ y)) \in r \vee (fst\ x = fst\ y \wedge ((snd\ x), (snd\ y)) \in s))$
 $\langle proof \rangle$

5.9 Context *od-term*

context *od-term*
begin

lemmas *red-wf = red-wf-dgrad-p-set*[*OF dickson-grading-zero subset-dgrad-p-set-zero*]
lemmas *Buchberger-criterion = Buchberger-criterion-dgrad-p-set*[*OF dickson-grading-zero subset-dgrad-p-set-zero*]

end

end

6 A General Algorithm Schema for Computing Gröbner Bases

theory *Algorithm-Schema*
imports *General Groebner-Bases*
begin

This theory formalizes a general algorithm schema for computing Gröbner bases, generalizing Buchberger's original critical-pair/completion algorithm.

The algorithm schema depends on several functional parameters that can be instantiated by a variety of concrete functions. Possible instances yield Buchberger's algorithm, Faugère's F4 algorithm, and (as far as we can tell) even his F5 algorithm.

6.1 *processed*

definition *minus-pairs* (**infixl** $-_p$ 65) **where** $\text{minus-pairs } A \ B = A - (B \cup \text{prod.swap } ' B)$

definition *Int-pairs* (**infixl** \cap_p 65) **where** $\text{Int-pairs } A \ B = A \cap (B \cup \text{prod.swap } ' B)$

definition *in-pair* (**infix** \in_p 50) **where** $\text{in-pair } p \ A \longleftrightarrow (p \in A \cup \text{prod.swap } ' A)$

definition *subset-pairs* (**infix** \subseteq_p 50) **where** $\text{subset-pairs } A \ B \longleftrightarrow (\forall x. x \in_p A \longrightarrow x \in_p B)$

abbreviation *not-in-pair* (**infix** \notin_p 50) **where** $\text{not-in-pair } p \ A \equiv \neg p \in_p A$

lemma *in-pair-alt*: $p \in_p A \longleftrightarrow (p \in A \vee \text{prod.swap } p \in A)$
<proof>

lemma *in-pair-iff*: $(a, b) \in_p A \longleftrightarrow ((a, b) \in A \vee (b, a) \in A)$
<proof>

lemma *in-pair-minus-pairs* [*simp*]: $p \in_p A -_p B \longleftrightarrow (p \in_p A \wedge p \notin_p B)$
<proof>

lemma *in-minus-pairs* [*simp*]: $p \in A -_p B \longleftrightarrow (p \in A \wedge p \notin_p B)$
<proof>

lemma *in-pair-Int-pairs* [*simp*]: $p \in_p A \cap_p B \longleftrightarrow (p \in_p A \wedge p \in_p B)$
<proof>

lemma *in-pair-Un* [*simp*]: $p \in_p A \cup B \longleftrightarrow (p \in_p A \vee p \in_p B)$
<proof>

lemma *in-pair-trans* [*trans*]:
assumes $p \in_p A$ **and** $A \subseteq B$
shows $p \in_p B$
<proof>

lemma *in-pair-same* [*simp*]: $p \in_p A \times A \longleftrightarrow p \in A \times A$
<proof>

lemma *subset-pairsI* [*intro*]:
assumes $\bigwedge x. x \in_p A \implies x \in_p B$
shows $A \subseteq_p B$
<proof>

lemma *subset-pairsD* [*trans*]:
assumes $x \in_p A$ **and** $A \subseteq_p B$

shows $x \in_p B$
<proof>

definition *processed* :: ('a × 'a) ⇒ 'a list ⇒ ('a × 'a) list ⇒ bool
where *processed* p xs ps $\longleftrightarrow p \in \text{set } xs \times \text{set } xs \wedge p \notin_p \text{set } ps$

lemma *processed-alt*:
processed (a, b) xs ps $\longleftrightarrow ((a \in \text{set } xs) \wedge (b \in \text{set } xs) \wedge (a, b) \notin_p \text{set } ps)$
<proof>

lemma *processedI*:
assumes $a \in \text{set } xs$ **and** $b \in \text{set } xs$ **and** $(a, b) \notin_p \text{set } ps$
shows *processed* (a, b) xs ps
<proof>

lemma *processedD1*:
assumes *processed* (a, b) xs ps
shows $a \in \text{set } xs$
<proof>

lemma *processedD2*:
assumes *processed* (a, b) xs ps
shows $b \in \text{set } xs$
<proof>

lemma *processedD3*:
assumes *processed* (a, b) xs ps
shows $(a, b) \notin_p \text{set } ps$
<proof>

lemma *processed-Nil*: *processed* (a, b) xs [] $\longleftrightarrow (a \in \text{set } xs \wedge b \in \text{set } xs)$
<proof>

lemma *processed-Cons*:
assumes *processed* (a, b) xs ps
and $a1: a = p \implies b = q \implies \text{thesis}$
and $a2: a = q \implies b = p \implies \text{thesis}$
and $a3: \text{processed } (a, b) \text{ xs } ((p, q) \# ps) \implies \text{thesis}$
shows *thesis*
<proof>

lemma *processed-minus*:
assumes *processed* (a, b) xs (ps -- qs)
and $a1: (a, b) \in_p \text{set } qs \implies \text{thesis}$
and $a2: \text{processed } (a, b) \text{ xs } ps \implies \text{thesis}$
shows *thesis*
<proof>

6.2 Algorithm Schema

6.2.1 *const-lt-component*

context *ordered-term*
begin

definition *const-lt-component* :: ($'t \Rightarrow_0 'b::zero$) \Rightarrow $'k$ option
where *const-lt-component* $p =$
 ($\text{let } v = \text{lt } p \text{ in if pp-of-term } v = 0 \text{ then Some (component-of-term } v) \text{ else None}$)

lemma *const-lt-component-SomeI*:
assumes $lp\ p = 0$ **and** $\text{component-of-term (lt } p) = \text{cmp}$
shows $\text{const-lt-component } p = \text{Some cmp}$
<proof>

lemma *const-lt-component-SomeD1*:
assumes $\text{const-lt-component } p = \text{Some cmp}$
shows $lp\ p = 0$
<proof>

lemma *const-lt-component-SomeD2*:
assumes $\text{const-lt-component } p = \text{Some cmp}$
shows $\text{component-of-term (lt } p) = \text{cmp}$
<proof>

lemma *const-lt-component-subset*:
 $\text{const-lt-component } ' (B - \{0\}) - \{None\} \subseteq \text{Some } ' \text{ component-of-term } ' \text{ Keys } B$
<proof>

corollary *card-const-lt-component-le*:
assumes *finite B*
shows $\text{card (const-lt-component } ' (B - \{0\}) - \{None\}) \leq \text{card (component-of-term } ' \text{ Keys } B)$
<proof>

end

6.2.2 Type synonyms

type-synonym ($'a, 'b, 'c$) *pdata'* = ($'a \Rightarrow_0 'b$) \times $'c$
type-synonym ($'a, 'b, 'c$) *pdata* = ($'a \Rightarrow_0 'b$) \times $\text{nat} \times 'c$
type-synonym ($'a, 'b, 'c$) *pdata-pair* = ($'a, 'b, 'c$) *pdata* \times ($'a, 'b, 'c$) *pdata*
type-synonym ($'a, 'b, 'c, 'd$) *selT* = ($'a, 'b, 'c$) *pdata list* \Rightarrow ($'a, 'b, 'c$) *pdata list*
 \Rightarrow
 ($'a, 'b, 'c$) *pdata-pair list* \Rightarrow $\text{nat} \times 'd \Rightarrow$ ($'a, 'b, 'c$)
pdata-pair list
type-synonym ($'a, 'b, 'c, 'd$) *complT* = ($'a, 'b, 'c$) *pdata list* \Rightarrow ($'a, 'b, 'c$) *pdata*

$list \Rightarrow$
 $('a, 'b, 'c) pdata-pair list \Rightarrow ('a, 'b, 'c) pdata-pair list$
 \Rightarrow
 $nat \times 'd \Rightarrow (('a, 'b, 'c) pdata' list \times 'd)$
type-synonym $('a, 'b, 'c, 'd) apT = ('a, 'b, 'c) pdata list \Rightarrow ('a, 'b, 'c) pdata list$
 \Rightarrow
 $('a, 'b, 'c) pdata-pair list \Rightarrow ('a, 'b, 'c) pdata list \Rightarrow$
 $nat \times 'd \Rightarrow$
 $('a, 'b, 'c) pdata-pair list$
type-synonym $('a, 'b, 'c, 'd) abT = ('a, 'b, 'c) pdata list \Rightarrow ('a, 'b, 'c) pdata list$
 \Rightarrow
 $('a, 'b, 'c) pdata list \Rightarrow nat \times 'd \Rightarrow ('a, 'b, 'c) pdata list$

6.2.3 Specification of the *selector* parameter

definition $sel-spec :: ('a, 'b, 'c, 'd) selT \Rightarrow bool$
where $sel-spec sel \longleftrightarrow$
 $(\forall gs bs ps data. ps \neq [] \longrightarrow (sel gs bs ps data \neq [] \wedge set (sel gs bs ps data)$
 $\subseteq set ps))$

lemma $sel-specI$:
assumes $\bigwedge gs bs ps data. ps \neq [] \Longrightarrow (sel gs bs ps data \neq [] \wedge set (sel gs bs ps data) \subseteq set ps)$
shows $sel-spec sel$
 $\langle proof \rangle$

lemma $sel-specD1$:
assumes $sel-spec sel$ **and** $ps \neq []$
shows $sel gs bs ps data \neq []$
 $\langle proof \rangle$

lemma $sel-specD2$:
assumes $sel-spec sel$ **and** $ps \neq []$
shows $set (sel gs bs ps data) \subseteq set ps$
 $\langle proof \rangle$

6.2.4 Specification of the *add-basis* parameter

definition $ab-spec :: ('a, 'b, 'c, 'd) abT \Rightarrow bool$
where $ab-spec ab \longleftrightarrow$
 $(\forall gs bs ns data. ns \neq [] \longrightarrow set (ab gs bs ns data) = set bs \cup set ns) \wedge$
 $(\forall gs bs data. ab gs bs [] data = bs)$

lemma $ab-specI$:
assumes $\bigwedge gs bs ns data. ns \neq [] \Longrightarrow set (ab gs bs ns data) = set bs \cup set ns$
and $\bigwedge gs bs data. ab gs bs [] data = bs$
shows $ab-spec ab$
 $\langle proof \rangle$

lemma $ab-specD1$:

assumes *ab-spec ab*
shows $set (ab\ gs\ bs\ ns\ data) = set\ bs \cup set\ ns$
 $\langle proof \rangle$

lemma *ab-specD2*:
assumes *ab-spec ab*
shows $ab\ gs\ bs \ []\ data = bs$
 $\langle proof \rangle$

6.2.5 Specification of the *add-pairs* parameter

definition *unique-idx* :: $('t, 'b, 'c) pdata\ list \Rightarrow (nat \times 'd) \Rightarrow bool$
where *unique-idx bs data* \longleftrightarrow
 $(\forall f \in set\ bs. \forall g \in set\ bs. fst\ (snd\ f) = fst\ (snd\ g) \longrightarrow f = g) \wedge$
 $(\forall f \in set\ bs. fst\ (snd\ f) < fst\ data)$

lemma *unique-idxI*:
assumes $\bigwedge f\ g. f \in set\ bs \Longrightarrow g \in set\ bs \Longrightarrow fst\ (snd\ f) = fst\ (snd\ g) \Longrightarrow f = g$
and $\bigwedge f. f \in set\ bs \Longrightarrow fst\ (snd\ f) < fst\ data$
shows *unique-idx bs data*
 $\langle proof \rangle$

lemma *unique-idxD1*:
assumes *unique-idx bs data* **and** $f \in set\ bs$ **and** $g \in set\ bs$ **and** $fst\ (snd\ f) = fst\ (snd\ g)$
 $(snd\ f) = (snd\ g)$
shows $f = g$
 $\langle proof \rangle$

lemma *unique-idxD2*:
assumes *unique-idx bs data* **and** $f \in set\ bs$
shows $fst\ (snd\ f) < fst\ data$
 $\langle proof \rangle$

lemma *unique-idx-Nil*: *unique-idx [] data*
 $\langle proof \rangle$

lemma *unique-idx-subset*:
assumes *unique-idx bs data* **and** $set\ bs' \subseteq set\ bs$
shows *unique-idx bs' data*
 $\langle proof \rangle$

context *gd-term*
begin

definition *ap-spec* :: $('t, 'b::field, 'c, 'd) apT \Rightarrow bool$
where *ap-spec ap* $\longleftrightarrow (\forall gs\ bs\ ps\ hs\ data. set\ (ap\ gs\ bs\ ps\ hs\ data) \subseteq set\ ps \cup (set\ hs \times (set\ gs \cup set\ bs \cup set\ hs)) \wedge$
 $(\forall B\ d\ m. \forall h \in set\ hs. \forall g \in set\ gs \cup set\ bs \cup set\ hs. dickson-grading\ d \longrightarrow$
 $set\ gs \cup set\ bs \cup set\ hs \subseteq B \longrightarrow fst\ ' B \subseteq dgrad-p-set\ d\ m \longrightarrow$

$$\begin{aligned}
& \text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs) \longrightarrow \text{unique-idx } (gs \text{ @ } bs \text{ @ } hs) \text{ data} \longrightarrow \\
& \text{is-Groebner-basis } (fst \text{ ' } \text{set } gs) \longrightarrow h \neq g \longrightarrow fst \ h \neq 0 \longrightarrow fst \ g \neq 0 \longrightarrow \\
& (\forall a \ b. (a, b) \in_p \text{set } (ap \ gs \ bs \ ps \ hs \ data) \longrightarrow fst \ a \neq 0 \longrightarrow fst \ b \neq 0 \longrightarrow \\
& \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ a) \ (fst \ b)) \longrightarrow \\
& (\forall a \ b. a \in \text{set } gs \cup \text{set } bs \longrightarrow b \in \text{set } gs \cup \text{set } bs \longrightarrow fst \ a \neq 0 \longrightarrow fst \ b \neq \\
0 \longrightarrow \\
& \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ a) \ (fst \ b)) \longrightarrow \\
& \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ h) \ (fst \ g)) \wedge \\
& (\forall B \ d \ m. \forall h \ g. \text{dickson-grading } d \longrightarrow \\
& \quad \text{set } gs \cup \text{set } bs \cup \text{set } hs \subseteq B \longrightarrow fst \text{ ' } B \subseteq \text{dgrad-p-set } d \ m \longrightarrow \\
& \quad \text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs) \longrightarrow (\text{set } gs \cup \text{set } bs) \cap \text{set } hs = \{\} \longrightarrow \\
& \quad \text{unique-idx } (gs \text{ @ } bs \text{ @ } hs) \text{ data} \longrightarrow \text{is-Groebner-basis } (fst \text{ ' } \text{set } gs) \longrightarrow \\
& \quad h \neq g \longrightarrow fst \ h \neq 0 \longrightarrow fst \ g \neq 0 \longrightarrow \\
& \quad (h, g) \in \text{set } ps \text{ -}_p \text{set } (ap \ gs \ bs \ ps \ hs \ data) \longrightarrow \\
& \quad (\forall a \ b. (a, b) \in_p \text{set } (ap \ gs \ bs \ ps \ hs \ data) \longrightarrow (a, b) \in_p \text{set } hs \times (\text{set } gs \cup \\
\text{set } bs \cup \text{set } hs) \longrightarrow \\
& \quad \quad \text{fst } a \neq 0 \longrightarrow \text{fst } b \neq 0 \longrightarrow \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ a) \\
(fst \ b)) \longrightarrow \\
& \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ h) \ (fst \ g)))
\end{aligned}$$

Informally, *ap-spec* *ap* means that, for suitable arguments *gs*, *bs*, *ps* and *hs*, the value of *ap gs bs ps hs* is a list of pairs *ps'* such that for every element (a, b) missing in *ps'* there exists a set of pairs *C* by reference to which (a, b) can be discarded, i. e. as soon as all critical pairs of the elements in *C* can be connected below some set *B*, the same is true for the critical pair of (a, b) .

lemma *ap-specI*:

$$\begin{aligned}
& \text{assumes } \bigwedge gs \ bs \ ps \ hs \ data. \text{set } (ap \ gs \ bs \ ps \ hs \ data) \subseteq \text{set } ps \cup (\text{set } hs \times (\text{set} \\
& \text{gs} \cup \text{set } bs \cup \text{set } hs)) \\
& \text{assumes } \bigwedge gs \ bs \ ps \ hs \ data \ B \ d \ m \ h \ g. \text{dickson-grading } d \Longrightarrow \\
& \quad \text{set } gs \cup \text{set } bs \cup \text{set } hs \subseteq B \Longrightarrow fst \text{ ' } B \subseteq \text{dgrad-p-set } d \ m \Longrightarrow \\
& \quad h \in \text{set } hs \Longrightarrow g \in \text{set } gs \cup \text{set } bs \cup \text{set } hs \Longrightarrow \\
& \quad \text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs) \Longrightarrow \text{unique-idx } (gs \text{ @ } bs \text{ @ } hs) \text{ data} \\
\Longrightarrow \\
& \quad \text{is-Groebner-basis } (fst \text{ ' } \text{set } gs) \Longrightarrow h \neq g \Longrightarrow fst \ h \neq 0 \Longrightarrow fst \ g \neq 0 \\
\Longrightarrow \\
& \quad (\bigwedge a \ b. (a, b) \in_p \text{set } (ap \ gs \ bs \ ps \ hs \ data) \Longrightarrow fst \ a \neq 0 \Longrightarrow fst \ b \neq 0 \\
\Longrightarrow \\
& \quad \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ a) \ (fst \ b)) \Longrightarrow \\
& \quad (\bigwedge a \ b. a \in \text{set } gs \cup \text{set } bs \Longrightarrow b \in \text{set } gs \cup \text{set } bs \Longrightarrow fst \ a \neq 0 \Longrightarrow \\
fst \ b \neq 0 \Longrightarrow \\
& \quad \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ a) \ (fst \ b)) \Longrightarrow \\
& \quad \text{crit-pair-cbelow-on } d \ m \ (fst \text{ ' } B) \ (fst \ h) \ (fst \ g)) \\
& \text{assumes } \bigwedge gs \ bs \ ps \ hs \ data \ B \ d \ m \ h \ g. \text{dickson-grading } d \Longrightarrow \\
& \quad \text{set } gs \cup \text{set } bs \cup \text{set } hs \subseteq B \Longrightarrow fst \text{ ' } B \subseteq \text{dgrad-p-set } d \ m \Longrightarrow \\
& \quad \text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs) \Longrightarrow (\text{set } gs \cup \text{set } bs) \cap \text{set } hs = \{\} \\
\Longrightarrow \\
& \quad \text{unique-idx } (gs \text{ @ } bs \text{ @ } hs) \text{ data} \Longrightarrow \text{is-Groebner-basis } (fst \text{ ' } \text{set } gs) \Longrightarrow \\
h \neq g \Longrightarrow
\end{aligned}$$

$fst\ h \neq 0 \implies fst\ g \neq 0 \implies (h, g) \in set\ ps \text{ } \text{-}_p\ set\ (ap\ gs\ bs\ ps\ hs\ data)$
 \implies
 $(\bigwedge a\ b. (a, b) \in_p set\ (ap\ gs\ bs\ ps\ hs\ data) \implies (a, b) \in_p set\ hs \times (set\ gs \cup set\ bs \cup set\ hs)) \implies$
 $fst\ a \neq 0 \implies fst\ b \neq 0 \implies crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ a)\ (fst\ b) \implies$
 $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ h)\ (fst\ g)$
shows *ap-spec ap*
<proof>

lemma *ap-specD1:*

assumes *ap-spec ap*

shows $set\ (ap\ gs\ bs\ ps\ hs\ data) \subseteq set\ ps \cup (set\ hs \times (set\ gs \cup set\ bs \cup set\ hs))$
<proof>

lemma *ap-specD2:*

assumes *ap-spec ap* **and** *dickson-grading d* **and** $set\ gs \cup set\ bs \cup set\ hs \subseteq B$
and $fst\ 'B \subseteq dgrad\text{-}p\text{-}set\ d\ m$ **and** $(h, g) \in_p set\ hs \times (set\ gs \cup set\ bs \cup set\ hs)$
and $set\ ps \subseteq set\ bs \times (set\ gs \cup set\ bs)$ **and** *unique-idx (gs @ bs @ hs) data*
and *is-Groebner-basis (fst ' set gs)* **and** $h \neq g$ **and** $fst\ h \neq 0$ **and** $fst\ g \neq 0$
and $\bigwedge a\ b. (a, b) \in_p set\ (ap\ gs\ bs\ ps\ hs\ data) \implies fst\ a \neq 0 \implies fst\ b \neq 0 \implies$
 $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ a)\ (fst\ b)$
and $\bigwedge a\ b. a \in set\ gs \cup set\ bs \implies b \in set\ gs \cup set\ bs \implies fst\ a \neq 0 \implies fst\ b$
 $\neq 0 \implies$
 $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ a)\ (fst\ b)$
shows $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ h)\ (fst\ g)$
<proof>

lemma *ap-specD3:*

assumes *ap-spec ap* **and** *dickson-grading d* **and** $set\ gs \cup set\ bs \cup set\ hs \subseteq B$
and $fst\ 'B \subseteq dgrad\text{-}p\text{-}set\ d\ m$ **and** $set\ ps \subseteq set\ bs \times (set\ gs \cup set\ bs)$
and $(set\ gs \cup set\ bs) \cap set\ hs = \{\}$ **and** *unique-idx (gs @ bs @ hs) data*
and *is-Groebner-basis (fst ' set gs)* **and** $h \neq g$ **and** $fst\ h \neq 0$ **and** $fst\ g \neq 0$
and $(h, g) \in_p set\ ps \text{ } \text{-}_p\ set\ (ap\ gs\ bs\ ps\ hs\ data)$
and $\bigwedge a\ b. a \in set\ hs \implies b \in set\ gs \cup set\ bs \cup set\ hs \implies (a, b) \in_p set\ (ap\ gs$
 $bs\ ps\ hs\ data) \implies$
 $fst\ a \neq 0 \implies fst\ b \neq 0 \implies crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ a)$
 $(fst\ b)$
shows $crit\text{-}pair\text{-}cbelow\text{-}on\ d\ m\ (fst\ 'B)\ (fst\ h)\ (fst\ g)$
<proof>

lemma *ap-spec-Nil-subset:*

assumes *ap-spec ap*

shows $set\ (ap\ gs\ bs\ ps\ \square\ data) \subseteq set\ ps$
<proof>

lemma *ap-spec-fst-subset:*

assumes *ap-spec ap*

shows $fst\ 'set\ (ap\ gs\ bs\ ps\ hs\ data) \subseteq fst\ 'set\ ps \cup set\ hs$

<proof>

lemma *ap-spec-snd-subset*:

assumes *ap-spec ap*

shows $\text{snd} \text{ ' set } (ap \text{ gs } bs \text{ ps } hs \text{ data}) \subseteq \text{snd} \text{ ' set } ps \cup \text{set } gs \cup \text{set } bs \cup \text{set } hs$

<proof>

lemma *ap-spec-inE*:

assumes *ap-spec ap* **and** $(p, q) \in \text{set } (ap \text{ gs } bs \text{ ps } hs \text{ data})$

assumes 1: $(p, q) \in \text{set } ps \implies \text{thesis}$

assumes 2: $p \in \text{set } hs \implies q \in \text{set } gs \cup \text{set } bs \cup \text{set } hs \implies \text{thesis}$

shows *thesis*

<proof>

lemma *subset-Times-ap*:

assumes *ap-spec ap* **and** *ab-spec ab* **and** $\text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs)$

shows $\text{set } (ap \text{ gs } bs \text{ (ps --- sps) } hs \text{ data}) \subseteq \text{set } (ab \text{ gs } bs \text{ hs } data) \times (\text{set } gs \cup \text{set } (ab \text{ gs } bs \text{ hs } data))$

<proof>

6.2.6 Function *args-to-set*

definition *args-to-set* :: $(t, 'b::\text{field}, 'c) \text{ pdata list} \times (t, 'b, 'c) \text{ pdata list} \times (t, 'b, 'c) \text{ pdata-pair list} \Rightarrow (t \Rightarrow_0 'b) \text{ set}$

where $\text{args-to-set } x = \text{fst} \text{ ' (set } (fst \text{ } x) \cup \text{set } (fst \text{ (snd } x)) \cup \text{fst} \text{ ' set } (snd \text{ (snd } x)) \cup \text{snd} \text{ ' set } (snd \text{ (snd } x)))$

lemma *args-to-set-alt*:

$\text{args-to-set } (gs, bs, ps) = \text{fst} \text{ ' set } gs \cup \text{fst} \text{ ' set } bs \cup \text{fst} \text{ ' fst} \text{ ' set } ps \cup \text{fst} \text{ ' snd} \text{ ' set } ps$

<proof>

lemma *args-to-set-subset-Times*:

assumes $\text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs)$

shows $\text{args-to-set } (gs, bs, ps) = \text{fst} \text{ ' set } gs \cup \text{fst} \text{ ' set } bs$

<proof>

lemma *args-to-set-subset*:

assumes *ap-spec ap* **and** *ab-spec ab*

shows $\text{args-to-set } (gs, ab \text{ gs } bs \text{ hs } data, ap \text{ gs } bs \text{ ps } hs \text{ data}) \subseteq$

$\text{fst} \text{ ' (set } gs \cup \text{set } bs \cup \text{fst} \text{ ' set } ps \cup \text{snd} \text{ ' set } ps \cup \text{set } hs) \text{ (is } ?l \subseteq \text{fst} \text{ '}$

$?r)$

<proof>

lemma *args-to-set-alt2*:

assumes *ap-spec ap* **and** *ab-spec ab* **and** $\text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs)$

shows $\text{args-to-set } (gs, ab \text{ gs } bs \text{ hs } data, ap \text{ gs } bs \text{ (ps --- sps) } hs \text{ data}) =$

$\text{fst} \text{ ' (set } gs \cup \text{set } bs \cup \text{set } hs) \text{ (is } ?l = \text{fst} \text{ ' } ?r)$

<proof>

lemma *args-to-set-subset1*:
assumes $set\ gs1 \subseteq set\ gs2$
shows $args\text{-to}\text{-set}\ (gs1, bs, ps) \subseteq args\text{-to}\text{-set}\ (gs2, bs, ps)$
 $\langle proof \rangle$

lemma *args-to-set-subset2*:
assumes $set\ bs1 \subseteq set\ bs2$
shows $args\text{-to}\text{-set}\ (gs, bs1, ps) \subseteq args\text{-to}\text{-set}\ (gs, bs2, ps)$
 $\langle proof \rangle$

lemma *args-to-set-subset3*:
assumes $set\ ps1 \subseteq set\ ps2$
shows $args\text{-to}\text{-set}\ (gs, bs, ps1) \subseteq args\text{-to}\text{-set}\ (gs, bs, ps2)$
 $\langle proof \rangle$

6.2.7 Functions *count-const-lt-components*, *count-rem-comps* and *full-gb*

definition *rem-comps-spec* :: $(t, 'b::zero, 'c)\ pdata\ list \Rightarrow nat \times 'd \Rightarrow bool$
where $rem\text{-comps}\text{-spec}\ bs\ data \longleftrightarrow (card\ (component\text{-of}\text{-term}\ 'Keys\ (fst\ 'set\ bs)) =$

$$fst\ data + card\ (const\text{-lt}\text{-component}\ ' (fst\ 'set\ bs - \{0\}) - \{None\}))$$

definition *count-const-lt-components* :: $(t, 'b::zero, 'c)\ pdata'\ list \Rightarrow nat$
where $count\text{-const}\text{-lt}\text{-components}\ hs = length\ (remdups\ (filter\ (\lambda x. x \neq None)$
 $(map\ (const\text{-lt}\text{-component}\ o\ fst)\ hs)))$

definition *count-rem-components* :: $(t, 'b::zero, 'c)\ pdata'\ list \Rightarrow nat$
where $count\text{-rem}\text{-components}\ bs = length\ (remdups\ (map\ component\text{-of}\text{-term}$
 $(Keys\text{-to}\text{-list}\ (map\ fst\ bs)))) -$
 $count\text{-const}\text{-lt}\text{-components}\ [b \leftarrow bs . fst\ b \neq 0]$

lemma *count-const-lt-components-alt*:
 $count\text{-const}\text{-lt}\text{-components}\ hs = card\ (const\text{-lt}\text{-component}\ 'fst\ 'set\ hs - \{None\})$
 $\langle proof \rangle$

lemma *count-rem-components-alt*:
 $count\text{-rem}\text{-components}\ bs + card\ (const\text{-lt}\text{-component}\ ' (fst\ 'set\ bs - \{0\}) -$
 $\{None\}) =$
 $card\ (component\text{-of}\text{-term}\ 'Keys\ (fst\ 'set\ bs))$
 $\langle proof \rangle$

lemma *rem-comps-spec-count-rem-components*: $rem\text{-comps}\text{-spec}\ bs\ (count\text{-rem}\text{-components}\ bs, data)$
 $\langle proof \rangle$

definition *full-gb* :: $(t, 'b, 'c)\ pdata\ list \Rightarrow (t, 'b::zero\text{-neq}\text{-one}, 'c::default)\ pdata\ list$

where $full\text{-}gb\ bs = map\ (\lambda k. (monomial\ 1\ (term\text{-}of\text{-}pair\ (0, k)), 0, default))$
 $(remdups\ (map\ component\text{-}of\text{-}term\ (Keys\text{-}to\text{-}list\ (map\ fst\ bs))))$

lemma $fst\text{-}set\text{-}full\text{-}gb$:

$fst\ 'set\ (full\text{-}gb\ bs) = (\lambda v. monomial\ 1\ (term\text{-}of\text{-}pair\ (0, component\text{-}of\text{-}term\ v)))$
 $'Keys\ (fst\ 'set\ bs)$
 $\langle proof \rangle$

lemma $Keys\text{-}full\text{-}gb$:

$Keys\ (fst\ 'set\ (full\text{-}gb\ bs)) = (\lambda v. term\text{-}of\text{-}pair\ (0, component\text{-}of\text{-}term\ v))\ 'Keys$
 $(fst\ 'set\ bs)$
 $\langle proof \rangle$

lemma $pps\text{-}full\text{-}gb$: $pp\text{-}of\text{-}term\ 'Keys\ (fst\ 'set\ (full\text{-}gb\ bs)) \subseteq \{0\}$

$\langle proof \rangle$

lemma $components\text{-}full\text{-}gb$:

$component\text{-}of\text{-}term\ 'Keys\ (fst\ 'set\ (full\text{-}gb\ bs)) = component\text{-}of\text{-}term\ 'Keys\ (fst$
 $'set\ bs)$
 $\langle proof \rangle$

lemma $full\text{-}gb\text{-}is\text{-}full\text{-}pmdl$: $is\text{-}full\text{-}pmdl\ (fst\ 'set\ (full\text{-}gb\ bs))$

for $bs::('t, 'b::field, 'c::default)\ pdata\ list$

$\langle proof \rangle$

In fact, $is\text{-}full\text{-}pmdl\ (fst\ 'set\ (full\text{-}gb\ ?bs))$ also holds if $'b$ is no field.

lemma $full\text{-}gb\text{-}isGB$: $is\text{-}Groebner\text{-}basis\ (fst\ 'set\ (full\text{-}gb\ bs))$

$\langle proof \rangle$

6.2.8 Specification of the *completion* parameter

definition $compl\text{-}struct :: ('t, 'b::field, 'c, 'd)\ complT \Rightarrow bool$

where $compl\text{-}struct\ compl \longleftrightarrow$

$(\forall gs\ bs\ ps\ sps\ data. sps \neq [] \longrightarrow set\ sps \subseteq set\ ps \longrightarrow$

$(\forall d. dickson\text{-}grading\ d \longrightarrow$

$dgrad\text{-}p\text{-}set\text{-}le\ d\ (fst\ '(set\ (fst\ (compl\ gs\ bs\ (ps\ \text{---}\ sps)\ sps\ data))))))$

$(args\text{-}to\text{-}set\ (gs, bs, ps))) \wedge$

$component\text{-}of\text{-}term\ 'Keys\ (fst\ '(set\ (fst\ (compl\ gs\ bs\ (ps\ \text{---}\ sps)\ sps$

$data)))) \subseteq$

$component\text{-}of\text{-}term\ 'Keys\ (args\text{-}to\text{-}set\ (gs, bs, ps)) \wedge$

$0 \notin fst\ 'set\ (fst\ (compl\ gs\ bs\ (ps\ \text{---}\ sps)\ sps\ data)) \wedge$

$(\forall h \in set\ (fst\ (compl\ gs\ bs\ (ps\ \text{---}\ sps)\ sps\ data)). \forall b \in set\ gs \cup set\ bs.$

$fst\ b \neq 0 \longrightarrow \neg lt\ (fst\ b)\ add_{st}\ lt\ (fst\ h)))$

lemma $compl\text{-}structI$:

assumes $\bigwedge d\ gs\ bs\ ps\ sps\ data. dickson\text{-}grading\ d \Longrightarrow sps \neq [] \Longrightarrow set\ sps \subseteq set$
 $ps \Longrightarrow$

$dgrad\text{-}p\text{-}set\text{-}le\ d\ (fst\ '(set\ (fst\ (compl\ gs\ bs\ (ps\ \text{---}\ sps)\ sps\ data))))$

$(args\text{-}to\text{-}set\ (gs, bs, ps))$

assumes $\bigwedge gs\ bs\ ps\ sps\ data.\ sps \neq [] \implies set\ sps \subseteq set\ ps \implies$
component-of-term 'Keys (fst ' (set (fst (compl gs bs (ps -- sps) sps data)))) \subseteq
component-of-term 'Keys (args-to-set (gs, bs, ps))
assumes $\bigwedge gs\ bs\ ps\ sps\ data.\ sps \neq [] \implies set\ sps \subseteq set\ ps \implies 0 \notin fst\ 'set\ (fst$
(compl gs bs (ps -- sps) sps data))
assumes $\bigwedge gs\ bs\ ps\ sps\ h\ b\ data.\ sps \neq [] \implies set\ sps \subseteq set\ ps \implies h \in set\ (fst$
(compl gs bs (ps -- sps) sps data)) \implies
 $b \in set\ gs \cup set\ bs \implies fst\ b \neq 0 \implies \neg lt\ (fst\ b)\ adds_t\ lt\ (fst\ h)$
shows *compl-struct compl*
<proof>

lemma *compl-structD1:*

assumes *compl-struct compl and dickson-grading d and sps $\neq []$ and set sps \subseteq*
set ps
shows *dgrad-p-set-le d (fst ' (set (fst (compl gs bs (ps -- sps) sps data))))*
(args-to-set (gs, bs, ps))
<proof>

lemma *compl-structD2:*

assumes *compl-struct compl and sps $\neq []$ and set sps \subseteq set ps*
shows *component-of-term 'Keys (fst ' (set (fst (compl gs bs (ps -- sps) sps data))))* \subseteq
component-of-term 'Keys (args-to-set (gs, bs, ps))
<proof>

lemma *compl-structD3:*

assumes *compl-struct compl and sps $\neq []$ and set sps \subseteq set ps*
shows $0 \notin fst\ 'set\ (fst\ (compl\ gs\ bs\ (ps\ --\ sps)\ sps\ data))$
<proof>

lemma *compl-structD4:*

assumes *compl-struct compl and sps $\neq []$ and set sps \subseteq set ps*
and $h \in set\ (fst\ (compl\ gs\ bs\ (ps\ --\ sps)\ sps\ data))$ **and** $b \in set\ gs \cup set\ bs$
and $fst\ b \neq 0$
shows $\neg lt\ (fst\ b)\ adds_t\ lt\ (fst\ h)$
<proof>

definition *struct-spec* :: $(t, 'b::field, 'c, 'd)\ selT \Rightarrow (t, 'b, 'c, 'd)\ apT \Rightarrow (t, 'b,$
 $'c, 'd)\ abT \Rightarrow$

$$(t, 'b, 'c, 'd)\ complT \Rightarrow bool$$

where *struct-spec sel ap ab compl* $\iff (sel\ spec\ sel \wedge ap\ spec\ ap \wedge ab\ spec\ ab \wedge$
compl-struct compl)

lemma *struct-specI:*

assumes *sel-spec sel and ap-spec ap and ab-spec ab and compl-struct compl*
shows *struct-spec sel ap ab compl*
<proof>

lemma *struct-specD1*:
assumes *struct-spec sel ap ab compl*
shows *sel-spec sel*
 \langle *proof* \rangle

lemma *struct-specD2*:
assumes *struct-spec sel ap ab compl*
shows *ap-spec ap*
 \langle *proof* \rangle

lemma *struct-specD3*:
assumes *struct-spec sel ap ab compl*
shows *ab-spec ab*
 \langle *proof* \rangle

lemma *struct-specD4*:
assumes *struct-spec sel ap ab compl*
shows *compl-struct compl*
 \langle *proof* \rangle

lemmas *struct-specD = struct-specD1 struct-specD2 struct-specD3 struct-specD4*

definition *compl-pmdl* :: $(t, 'b::field, 'c, 'd)$ *complT* \Rightarrow *bool*
where *compl-pmdl compl* \longleftrightarrow
 $(\forall gs\ bs\ ps\ sps\ data. is-Groebner-basis\ (fst\ 'set\ gs) \longrightarrow sps \neq [] \longrightarrow set$
 $sps \subseteq set\ ps \longrightarrow$
 $unique-idx\ (gs\ @\ bs)\ data \longrightarrow$
 $fst\ '(\set\ (fst\ (compl\ gs\ bs\ (ps\ --\ sps)\ sps\ data))) \subseteq pmdl\ (args-to-set$
 $(gs,\ bs,\ ps)))$

lemma *compl-pmdlI*:
assumes $\bigwedge gs\ bs\ ps\ sps\ data. is-Groebner-basis\ (fst\ 'set\ gs) \Longrightarrow sps \neq [] \Longrightarrow set$
 $sps \subseteq set\ ps \Longrightarrow$
 $unique-idx\ (gs\ @\ bs)\ data \Longrightarrow$
 $fst\ '(\set\ (fst\ (compl\ gs\ bs\ (ps\ --\ sps)\ sps\ data))) \subseteq pmdl\ (args-to-set$
 $(gs,\ bs,\ ps))$
shows *compl-pmdl compl*
 \langle *proof* \rangle

lemma *compl-pmdlD*:
assumes *compl-pmdl compl and is-Groebner-basis (fst 'set gs)*
and $sps \neq []$ **and** $set\ sps \subseteq set\ ps$ **and** *unique-idx (gs @ bs) data*
shows $fst\ '(\set\ (fst\ (compl\ gs\ bs\ (ps\ --\ sps)\ sps\ data))) \subseteq pmdl\ (args-to-set$
 $(gs,\ bs,\ ps))$
 \langle *proof* \rangle

definition *compl-conn* :: $(t, 'b::field, 'c, 'd)$ *complT* \Rightarrow *bool*
where *compl-conn compl* \longleftrightarrow
 $(\forall d\ m\ gs\ bs\ ps\ sps\ p\ q\ data. dickson-grading\ d \longrightarrow fst\ 'set\ gs \subseteq dgrad-p-set$

$d \ m \longrightarrow$
 $is\text{-Groebner-basis} \ (fst \ ' \ set \ gs) \longrightarrow fst \ ' \ set \ bs \subseteq \ dgrad\text{-p-set} \ d \ m \longrightarrow$
 $set \ ps \subseteq set \ bs \times (set \ gs \cup set \ bs) \longrightarrow sps \neq [] \longrightarrow set \ sps \subseteq set \ ps \longrightarrow$
 $unique\text{-idx} \ (gs \ @ \ bs) \ data \longrightarrow (p, q) \in set \ sps \longrightarrow fst \ p \neq 0 \longrightarrow fst \ q$
 $\neq 0 \longrightarrow$
 $crit\text{-pair-cbelow-on} \ d \ m \ (fst \ ' \ (set \ gs \cup set \ bs) \cup fst \ ' \ set \ (fst \ (compl \ gs$
 $bs \ (ps \ \text{--} \ sps) \ sps \ data))) \ (fst \ p) \ (fst \ q)$

Informally, *compl-conn compl* means that, for suitable arguments *gs*, *bs*, *ps* and *sps*, the value of *compl gs bs ps sps* is a list *hs* such that the critical pairs of all elements in *sps* can be connected modulo *set gs* \cup *set bs* \cup *set hs*.

lemma *compl-connI*:

assumes $\bigwedge d \ m \ gs \ bs \ ps \ sps \ p \ q \ data. \ dickson\text{-grading} \ d \Longrightarrow fst \ ' \ set \ gs \subseteq$
 $dgrad\text{-p-set} \ d \ m \Longrightarrow$
 $is\text{-Groebner-basis} \ (fst \ ' \ set \ gs) \Longrightarrow fst \ ' \ set \ bs \subseteq \ dgrad\text{-p-set} \ d \ m \Longrightarrow$
 $set \ ps \subseteq set \ bs \times (set \ gs \cup set \ bs) \Longrightarrow sps \neq [] \Longrightarrow set \ sps \subseteq set \ ps \Longrightarrow$
 $unique\text{-idx} \ (gs \ @ \ bs) \ data \Longrightarrow (p, q) \in set \ sps \Longrightarrow fst \ p \neq 0 \Longrightarrow fst \ q \neq$
 $0 \Longrightarrow$
 $crit\text{-pair-cbelow-on} \ d \ m \ (fst \ ' \ (set \ gs \cup set \ bs) \cup fst \ ' \ set \ (fst \ (compl \ gs \ bs$
 $(ps \ \text{--} \ sps) \ sps \ data))) \ (fst \ p) \ (fst \ q)$
shows *compl-conn compl*
 $\langle proof \rangle$

lemma *compl-connD*:

assumes *compl-conn compl* **and** *dickson-grading d* **and** $fst \ ' \ set \ gs \subseteq \ dgrad\text{-p-set}$
 $d \ m$
and *is-Groebner-basis (fst ' set gs)* **and** $fst \ ' \ set \ bs \subseteq \ dgrad\text{-p-set} \ d \ m$
and $set \ ps \subseteq set \ bs \times (set \ gs \cup set \ bs)$ **and** $sps \neq []$ **and** $set \ sps \subseteq set \ ps$
and *unique-idx (gs @ bs) data* **and** $(p, q) \in set \ sps$ **and** $fst \ p \neq 0$ **and** $fst \ q \neq$
 0
shows *crit-pair-cbelow-on d m (fst ' (set gs* \cup *set bs) \cup fst ' set (fst (compl gs*
 $(ps \ \text{--} \ sps) \ sps \ data))) \ (fst \ p) \ (fst \ q)$
 $\langle proof \rangle$

6.2.9 Function *gb-schema-dummy*

definition (**in** $-$) *add-indices* $:: ((\ 'a, \ 'b, \ 'c) \ pdata' \ list \times \ 'd) \Rightarrow (nat \times \ 'd) \Rightarrow ((\ 'a,$
 $\ 'b, \ 'c) \ pdata \ list \times \ nat \times \ 'd)$

where $[code \ del]: \ add\text{-indices} \ ns \ data =$
 $(map\text{-idx} \ (\lambda h \ i. \ (fst \ h, \ i, \ snd \ h)) \ (fst \ ns) \ (fst \ data), \ fst \ data + length \ (fst$
 $ns), \ snd \ ns)$

lemma (**in** $-$) *add-indices-code* $[code]:$

$add\text{-indices} \ (ns, \ data) \ (n, \ data') = (map\text{-idx} \ (\lambda(h, \ d) \ i. \ (h, \ i, \ d)) \ ns \ n, \ n + length$
 $ns, \ data)$
 $\langle proof \rangle$

lemma *fst-add-indices*: $map \ fst \ (fst \ (add\text{-indices} \ ns \ data')) = map \ fst \ (fst \ ns)$

<proof>

corollary *fst-set-add-indices*: $\text{fst } \text{'set } (\text{fst } (\text{add-indices } ns \text{ data}')) = \text{fst } \text{'set } (\text{fst } ns)$

<proof>

lemma *in-set-add-indicesE*:

assumes $f \in \text{set } (\text{fst } (\text{add-indices } aux \text{ data}))$

obtains i **where** $i < \text{length } (\text{fst } aux)$ **and** $f = (\text{fst } ((\text{fst } aux) ! i), \text{fst } data + i, \text{snd } ((\text{fst } aux) ! i))$

<proof>

definition *gb-schema-aux-term1* :: $((('t, 'b::\text{field}, 'c) \text{pdata list} \times ('t, 'b, 'c) \text{pdata-pair list}) \times$

$((('t, 'b, 'c) \text{pdata list} \times ('t, 'b, 'c) \text{pdata-pair list})) \text{set}$

where $\text{gb-schema-aux-term1} = \{(a, b::('t, 'b, 'c) \text{pdata list}). (\text{fst } \text{'set } a) \sqsupseteq p (\text{fst } \text{'set } b)\}$ *< *lex* >*

$(\text{measure } (\text{card } \circ \text{set}))$

definition *gb-schema-aux-term2* ::

$(\text{'a} \Rightarrow \text{nat}) \Rightarrow ('t, 'b::\text{field}, 'c) \text{pdata list} \Rightarrow (((('t, 'b, 'c) \text{pdata list} \times ('t, 'b, 'c) \text{pdata-pair list}) \times$

$((('t, 'b, 'c) \text{pdata list} \times ('t, 'b, 'c) \text{pdata-pair list})) \text{set}$

where $\text{gb-schema-aux-term2 } d \text{ gs} = \{(a, b). \text{dgrad-p-set-le } d (\text{args-to-set } (gs, a)) (\text{args-to-set } (gs, b)) \wedge$

$\text{component-of-term } \text{'Keys } (\text{args-to-set } (gs, a)) \subseteq \text{component-of-term}$

$\text{'Keys } (\text{args-to-set } (gs, b))\}$

definition *gb-schema-aux-term* **where** $\text{gb-schema-aux-term } d \text{ gs} = \text{gb-schema-aux-term1} \cap \text{gb-schema-aux-term2 } d \text{ gs}$

gb-schema-aux-term is needed for proving termination of function *gb-schema-aux*.

lemma *gb-schema-aux-term1-wf-on*:

assumes *dickson-grading* d **and** *finite* K

shows *wfp-on* $(\lambda x y. (x, y) \in \text{gb-schema-aux-term1})$

$\{x::(('t, 'b, 'c) \text{pdata list}) \times (((('t, 'b::\text{field}, 'c) \text{pdata-pair list})).$

$\text{args-to-set } (gs, x) \subseteq \text{dgrad-p-set } d \text{ m} \wedge \text{component-of-term } \text{'Keys}$

$(\text{args-to-set } (gs, x)) \subseteq K\}$

<proof>

lemma *gb-schema-aux-term-wf*:

assumes *dickson-grading* d

shows *wf* $(\text{gb-schema-aux-term } d \text{ gs})$

<proof>

lemma *dgrad-p-set-le-args-to-set-ab*:

assumes *dickson-grading* d **and** *ap-spec* ap **and** *ab-spec* ab **and** *compl-struct* compl

assumes $\text{sps} \neq []$ **and** $\text{set } \text{sps} \subseteq \text{set } \text{ps}$ **and** $\text{hs} = \text{fst } (\text{add-indices } (\text{compl } \text{gs } \text{bs}$

($ps \dashv\vdash sps$) sps data) data)
shows $dgrad-p-set-le$ d ($args-to-set$ (gs , ab gs bs hs data'), ap gs bs ($ps \dashv\vdash sps$)
 hs data')) ($args-to-set$ (gs , bs , ps))
(is $dgrad-p-set-le$ - ? l ? r)
⟨proof⟩

corollary $dgrad-p-set-le-args-to-set-struct$:

assumes $dickson-grading$ d **and** $struct-spec$ sel ap ab $compl$ **and** $ps \neq \square$
assumes $sps = sel$ gs bs ps data **and** $hs = fst$ ($add-indices$ ($compl$ gs bs ($ps \dashv\vdash$
 sps) sps data) data)
shows $dgrad-p-set-le$ d ($args-to-set$ (gs , ab gs bs hs data'), ap gs bs ($ps \dashv\vdash$
 hs data')) ($args-to-set$ (gs , bs , ps))
⟨proof⟩

lemma $components-subset-ab$:

assumes $ap-spec$ ap **and** $ab-spec$ ab **and** $compl-struct$ $compl$
assumes $sps \neq \square$ **and** set $sps \subseteq set$ ps **and** $hs = fst$ ($add-indices$ ($compl$ gs bs
($ps \dashv\vdash sps$) sps data) data)
shows $component-of-term$ ' $Keys$ ($args-to-set$ (gs , ab gs bs hs data'), ap gs bs (ps
 $\dashv\vdash sps$) hs data')) \subseteq
 $component-of-term$ ' $Keys$ ($args-to-set$ (gs , bs , ps)) (is ? $l \subseteq$? r)
⟨proof⟩

corollary $components-subset-struct$:

assumes $struct-spec$ sel ap ab $compl$ **and** $ps \neq \square$
assumes $sps = sel$ gs bs ps data **and** $hs = fst$ ($add-indices$ ($compl$ gs bs ($ps \dashv\vdash$
 sps) sps data) data)
shows $component-of-term$ ' $Keys$ ($args-to-set$ (gs , ab gs bs hs data'), ap gs bs (ps
 $\dashv\vdash sps$) hs data')) \subseteq
 $component-of-term$ ' $Keys$ ($args-to-set$ (gs , bs , ps))
⟨proof⟩

corollary $components-struct$:

assumes $struct-spec$ sel ap ab $compl$ **and** $ps \neq \square$ **and** set $ps \subseteq set$ $bs \times (set$ gs
 $\cup set$ $bs)$
assumes $sps = sel$ gs bs ps data **and** $hs = fst$ ($add-indices$ ($compl$ gs bs ($ps \dashv\vdash$
 sps) sps data) data)
shows $component-of-term$ ' $Keys$ ($args-to-set$ (gs , ab gs bs hs data'), ap gs bs (ps
 $\dashv\vdash sps$) hs data')) =
 $component-of-term$ ' $Keys$ ($args-to-set$ (gs , bs , ps)) (is ? $l =$? r)
⟨proof⟩

lemma $struct-spec-red-supset$:

assumes $struct-spec$ sel ap ab $compl$ **and** $ps \neq \square$ **and** $sps = sel$ gs bs ps data
and $hs = fst$ ($add-indices$ ($compl$ gs bs ($ps \dashv\vdash sps$) sps data) data) **and** $hs \neq$
 \square
shows (fst ' set (ab gs bs hs data')) \sqsupseteq (fst ' set bs)
⟨proof⟩

lemma *unique-idx-append*:

assumes *unique-idx gs data* **and** $(hs, data') = \text{add-indices aux data}$

shows *unique-idx (gs @ hs) data'*

<proof>

corollary *unique-idx-ab*:

assumes *ab-spec ab* **and** *unique-idx (gs @ bs) data* **and** $(hs, data') = \text{add-indices aux data}$

shows *unique-idx (gs @ ab gs bs hs data') data'*

<proof>

lemma *rem-comps-spec-struct*:

assumes *struct-spec sel ap ab compl* **and** *rem-comps-spec (gs @ bs) data* **and** $ps \neq \square$

and $set\ ps \subseteq (set\ bs) \times (set\ gs \cup set\ bs)$ **and** $sps = sel\ gs\ bs\ ps\ (snd\ data)$

and $aux = compl\ gs\ bs\ (ps \dashv\ dashv\ sps)\ sps\ (snd\ data)$ **and** $(hs, data') = \text{add-indices aux (snd data)}$

shows *rem-comps-spec (gs @ ab gs bs hs data')* $(fst\ data - \text{count-const-lt-components } (fst\ aux), data')$

<proof>

lemma *pmdl-struct*:

assumes *struct-spec sel ap ab compl* **and** *compl-pmdl compl* **and** *is-Groebner-basis (fst ' set gs)*

and $ps \neq \square$ **and** $set\ ps \subseteq (set\ bs) \times (set\ gs \cup set\ bs)$ **and** *unique-idx (gs @ bs) (snd data)*

and $sps = sel\ gs\ bs\ ps\ (snd\ data)$ **and** $aux = compl\ gs\ bs\ (ps \dashv\ dashv\ sps)\ sps\ (snd\ data)$

and $(hs, data') = \text{add-indices aux (snd data)}$

shows $pmdl\ (fst\ ' \ set\ (gs\ @\ ab\ gs\ bs\ hs\ data')) = pmdl\ (fst\ ' \ set\ (gs\ @\ bs))$

<proof>

lemma *discarded-subset*:

assumes *ab-spec ab*

and $D' = D \cup (set\ hs \times (set\ gs \cup set\ bs \cup set\ hs) \cup set\ (ps \dashv\ dashv\ sps) \dashv\ dashv\ set\ (ap\ gs\ bs\ (ps \dashv\ dashv\ sps)\ hs\ data'))$

and $set\ ps \subseteq set\ bs \times (set\ gs \cup set\ bs)$ **and** $D \subseteq (set\ gs \cup set\ bs) \times (set\ gs \cup set\ bs)$

shows $D' \subseteq (set\ gs \cup set\ (ab\ gs\ bs\ hs\ data')) \times (set\ gs \cup set\ (ab\ gs\ bs\ hs\ data'))$

<proof>

lemma *compl-struct-disjoint*:

assumes *compl-struct compl* **and** $sps \neq \square$ **and** $set\ sps \subseteq set\ ps$

shows $fst\ ' \ set\ (fst\ (compl\ gs\ bs\ (ps \dashv\ dashv\ sps)\ sps\ data)) \cap fst\ ' \ (set\ gs \cup set\ bs) = \{\}$

<proof>

context

$remcomps = fst (data) - count-const-lt-components (fst aux)$ in
 (if $remcomps = 0$ then
 $(full-gb (gs @ bs), D)$
 else
 let $(hs, data') = add-indices aux (snd data)$ in
 $gb-schema-dummy (remcomps, data')$
 $(D \cup ((set hs \times (set gs \cup set bs \cup set hs) \cup set (ps -- sps))) -_p$
 $set (ap gs bs ps0 hs data'))$
 $(ab gs bs hs data') (ap gs bs ps0 hs data')$
)
)
 <proof>

lemma *gb-schema-dummy-induct* [consumes 1, case-names base rec1 rec2]:

assumes *struct-spec sel ap ab compl*
assumes base: $\bigwedge bs data D. P data D bs [] (gs @ bs, D)$
and rec1: $\bigwedge bs ps sps data D. ps \neq [] \implies sps = sel gs bs ps (snd data) \implies$
 $fst (data) \leq count-const-lt-components (fst (compl gs bs (ps -- sps)$
 $sps (snd data))) \implies$
 $P data D bs ps (full-gb (gs @ bs), D)$
and rec2: $\bigwedge bs ps sps aux hs rc data data' D D'. ps \neq [] \implies sps = sel gs bs ps$
 $(snd data) \implies$
 $aux = compl gs bs (ps -- sps) sps (snd data) \implies (hs, data') =$
 $add-indices aux (snd data) \implies$
 $rc = fst data - count-const-lt-components (fst aux) \implies 0 < rc \implies$
 $D' = (D \cup ((set hs \times (set gs \cup set bs \cup set hs) \cup set (ps -- sps)))$
 $-_p set (ap gs bs (ps -- sps) hs data')) \implies$
 $P (rc, data') D' (ab gs bs hs data') (ap gs bs (ps -- sps) hs data')$
 $(gb-schema-dummy (rc, data') D' (ab gs bs hs data') (ap gs bs (ps$
 $-- sps) hs data')) \implies$
 $P data D bs ps (gb-schema-dummy (rc, data') D' (ab gs bs hs data')$
 $(ap gs bs (ps -- sps) hs data'))$
shows $P data D bs ps (gb-schema-dummy data D bs ps)$
 <proof>

lemma *fst-gb-schema-dummy-dgrad-p-set-le*:

assumes *dickson-grading d and struct-spec sel ap ab compl*
shows *dgrad-p-set-le d (fst ' set (fst (gb-schema-dummy data D bs ps))) (args-to-set*
 $(gs, bs, ps))$
 <proof>

lemma *fst-gb-schema-dummy-components*:

assumes *struct-spec sel ap ab compl and set ps $\subseteq (set bs) \times (set gs \cup set bs)$*
shows *component-of-term ' Keys (fst ' set (fst (gb-schema-dummy data D bs ps)))*
 $=$
 $component-of-term ' Keys (args-to-set (gs, bs, ps))$
 <proof>

lemma *fst-gb-schema-dummy-pmdl*:

assumes *struct-spec sel ap ab compl and compl-pmdl compl and is-Groebner-basis*
(fst ' set gs)
and *set ps* \subseteq *set bs* \times (*set gs* \cup *set bs*) **and** *unique-idx (gs @ bs) (snd data)*
and *rem-comps-spec (gs @ bs) data*
shows *pmdl (fst ' set (fst (gb-schema-dummy data D bs ps)))* = *pmdl (fst ' set*
(gs @ bs))
<proof>

lemma *snd-gb-schema-dummy-subset:*

assumes *struct-spec sel ap ab compl and set ps* \subseteq *set bs* \times (*set gs* \cup *set bs*)
and *D* \subseteq (*set gs* \cup *set bs*) \times (*set gs* \cup *set bs*) **and** *res = gb-schema-dummy*
data D bs ps
shows *snd res* \subseteq *set (fst res)* \times *set (fst res)* \vee ($\exists xs.$ *fst (res) = full-gb xs*)
<proof>

lemma *gb-schema-dummy-connectible1:*

assumes *struct-spec sel ap ab compl and compl-conn compl and dickson-grading*
d
and *fst ' set gs* \subseteq *dgrad-p-set d m* **and** *is-Groebner-basis (fst ' set gs)*
and *fst ' set bs* \subseteq *dgrad-p-set d m*
and *set ps* \subseteq *set bs* \times (*set gs* \cup *set bs*)
and *unique-idx (gs @ bs) (snd data)*
and $\bigwedge p q.$ *processed (p, q) (gs @ bs) ps* \implies $(p, q) \notin_p D \implies$ *fst p* $\neq 0 \implies$ *fst*
q $\neq 0 \implies$
crit-pair-cbelow-on d m (fst ' (set gs \cup set bs)) (fst p) (fst q)
and $\neg(\exists xs.$ *fst (gb-schema-dummy data D bs ps) = full-gb xs*)
assumes *f* \in *set (fst (gb-schema-dummy data D bs ps))*
and *g* \in *set (fst (gb-schema-dummy data D bs ps))*
and $(f, g) \notin_p$ *snd (gb-schema-dummy data D bs ps)*
and *fst f* $\neq 0$ **and** *fst g* $\neq 0$
shows *crit-pair-cbelow-on d m (fst ' set (fst (gb-schema-dummy data D bs ps)))*
(fst f) (fst g)
<proof>

lemma *gb-schema-dummy-connectible2:*

assumes *struct-spec sel ap ab compl and compl-conn compl and dickson-grading*
d
and *fst ' set gs* \subseteq *dgrad-p-set d m* **and** *is-Groebner-basis (fst ' set gs)*
and *fst ' set bs* \subseteq *dgrad-p-set d m*
and *set ps* \subseteq *set bs* \times (*set gs* \cup *set bs*) **and** *D* \subseteq (*set gs* \cup *set bs*) \times (*set gs* \cup
set bs)
and *set ps* $\cap_p D = \{\}$ **and** *unique-idx (gs @ bs) (snd data)*
and $\bigwedge B a b.$ *set gs* \cup *set bs* $\subseteq B \implies$ *fst ' B* \subseteq *dgrad-p-set d m* \implies $(a, b) \in_p$
D \implies
fst a $\neq 0 \implies$ *fst b* $\neq 0 \implies$
 $(\bigwedge x y. x \in$ *set gs* \cup *set bs* $\implies y \in$ *set gs* \cup *set bs* $\implies \neg(x, y) \in_p D \implies$
fst x $\neq 0 \implies$ *fst y* $\neq 0 \implies$ *crit-pair-cbelow-on d m (fst ' B) (fst x)*
(fst y)) \implies
crit-pair-cbelow-on d m (fst ' B) (fst a) (fst b)

and $\bigwedge x y. x \in \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps})) \implies y \in \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps})) \implies$
 $(x, y) \notin_p \text{snd } (\text{gb-schema-dummy data } D \text{ bs ps}) \implies \text{fst } x \neq 0 \implies \text{fst } y \neq 0 \implies$
 $\text{crit-pair-cbelow-on } d \text{ m } (\text{fst } ' \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps})))$
 $(\text{fst } x) (\text{fst } y)$
and $\neg(\exists xs. \text{fst } (\text{gb-schema-dummy data } D \text{ bs ps}) = \text{full-gb } xs)$
assumes $(f, g) \in_p \text{snd } (\text{gb-schema-dummy data } D \text{ bs ps})$
and $\text{fst } f \neq 0$ **and** $\text{fst } g \neq 0$
shows $\text{crit-pair-cbelow-on } d \text{ m } (\text{fst } ' \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps})))$
 $(\text{fst } f) (\text{fst } g)$
 $\langle \text{proof} \rangle$

corollary *gb-schema-dummy-connectible:*

assumes *struct-spec sel ap ab compl* **and** *compl-conn compl* **and** *dickson-grading*
 d
and $\text{fst } ' \text{set } gs \subseteq \text{dgrad-p-set } d \text{ m}$ **and** *is-Groebner-basis* $(\text{fst } ' \text{set } gs)$
and $\text{fst } ' \text{set } bs \subseteq \text{dgrad-p-set } d \text{ m}$
and $\text{set } ps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs)$ **and** $D \subseteq (\text{set } gs \cup \text{set } bs) \times (\text{set } gs \cup \text{set } bs)$
and $\text{set } ps \cap_p D = \{\}$ **and** *unique-idx* $(gs @ bs) (\text{snd } \text{data})$
and $\bigwedge p q. \text{processed } (p, q) (gs @ bs) ps \implies (p, q) \notin_p D \implies \text{fst } p \neq 0 \implies \text{fst } q \neq 0 \implies$
 $\text{crit-pair-cbelow-on } d \text{ m } (\text{fst } ' (\text{set } gs \cup \text{set } bs)) (\text{fst } p) (\text{fst } q)$
and $\bigwedge B a b. \text{set } gs \cup \text{set } bs \subseteq B \implies \text{fst } ' B \subseteq \text{dgrad-p-set } d \text{ m} \implies (a, b) \in_p$
 $D \implies$
 $\text{fst } a \neq 0 \implies \text{fst } b \neq 0 \implies$
 $(\bigwedge x y. x \in \text{set } gs \cup \text{set } bs \implies y \in \text{set } gs \cup \text{set } bs \implies \neg(x, y) \in_p D \implies$
 $\text{fst } x \neq 0 \implies \text{fst } y \neq 0 \implies \text{crit-pair-cbelow-on } d \text{ m } (\text{fst } ' B) (\text{fst } x)$
 $(\text{fst } y)) \implies$
 $\text{crit-pair-cbelow-on } d \text{ m } (\text{fst } ' B) (\text{fst } a) (\text{fst } b)$
assumes $f \in \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps}))$
and $g \in \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps}))$
and $\text{fst } f \neq 0$ **and** $\text{fst } g \neq 0$
shows $\text{crit-pair-cbelow-on } d \text{ m } (\text{fst } ' \text{set } (\text{fst } (\text{gb-schema-dummy data } D \text{ bs ps})))$
 $(\text{fst } f) (\text{fst } g)$
 $\langle \text{proof} \rangle$

lemma *fst-gb-schema-dummy-dgrad-p-set-le-init:*

assumes *dickson-grading* d **and** *struct-spec sel ap ab compl*
shows $\text{dgrad-p-set-le } d (\text{fst } ' \text{set } (\text{fst } (\text{gb-schema-dummy data } D (\text{ab } gs \square bs (\text{snd } \text{data}))) (\text{ap } gs \square \square bs (\text{snd } \text{data}))))))$
 $(\text{fst } ' (\text{set } gs \cup \text{set } bs))$
 $\langle \text{proof} \rangle$

corollary *fst-gb-schema-dummy-dgrad-p-set-init:*

assumes *dickson-grading* d **and** *struct-spec sel ap ab compl*
and $\text{fst } ' (\text{set } gs \cup \text{set } bs) \subseteq \text{dgrad-p-set } d \text{ m}$
shows $\text{fst } ' \text{set } (\text{fst } (\text{gb-schema-dummy } (rc, \text{data}) D (\text{ab } gs \square bs \text{data}))) (\text{ap } gs \square \square$

$bs\ data))) \subseteq dgrad-p-set\ d\ m$
 $\langle proof \rangle$

lemma *fst-gb-schema-dummy-components-init:*

fixes $bs\ data$
defines $bs0 \equiv ab\ gs \ []\ bs\ data$
defines $ps0 \equiv ap\ gs \ []\ []\ bs\ data$
assumes *struct-spec sel ap ab compl*
shows $component-of-term\ 'Keys\ (fst\ 'set\ (fst\ (gb-schema-dummy\ (rc,\ data)\ D\ bs0\ ps0))) =$
 $component-of-term\ 'Keys\ (fst\ 'set\ (gs\ @\ bs))\ (is\ ?l = ?r)$
 $\langle proof \rangle$

lemma *fst-gb-schema-dummy-pmdl-init:*

fixes $bs\ data$
defines $bs0 \equiv ab\ gs \ []\ bs\ data$
defines $ps0 \equiv ap\ gs \ []\ []\ bs\ data$
assumes *struct-spec sel ap ab compl and compl-pmdl compl and is-Groebner-basis*
 $(fst\ 'set\ gs)$
and *unique-idx (gs @ bs0) data and rem-comps-spec (gs @ bs0) (rc, data)*
shows $pmdl\ (fst\ 'set\ (fst\ (gb-schema-dummy\ (rc,\ data)\ D\ bs0\ ps0))) =$
 $pmdl\ (fst\ '(set\ (gs\ @\ bs)))\ (is\ ?l = ?r)$
 $\langle proof \rangle$

lemma *fst-gb-schema-dummy-isGB-init:*

fixes $bs\ data$
defines $bs0 \equiv ab\ gs \ []\ bs\ data$
defines $ps0 \equiv ap\ gs \ []\ []\ bs\ data$
defines $D0 \equiv set\ bs \times (set\ gs \cup set\ bs) \ -_p\ set\ ps0$
assumes *struct-spec sel ap ab compl and compl-conn compl and is-Groebner-basis*
 $(fst\ 'set\ gs)$
and *unique-idx (gs @ bs0) data and rem-comps-spec (gs @ bs0) (rc, data)*
shows *is-Groebner-basis (fst 'set (fst (gb-schema-dummy (rc, data) D0 bs0 ps0)))*
 $\langle proof \rangle$

6.2.10 Function *gb-schema-aux*

function $(domintros)\ gb-schema-aux :: nat \times nat \times 'd \Rightarrow ('t, 'b, 'c)\ pdata\ list \Rightarrow$
 $('t, 'b, 'c)\ pdata-pair\ list \Rightarrow ('t, 'b, 'c)\ pdata\ list$

where

$gb-schema-aux\ data\ bs\ ps =$
 $(if\ ps = []\ then$
 $gs\ @\ bs$
 $else$
 $(let\ sps = sel\ gs\ bs\ ps\ (snd\ data); ps0 = ps \ -\ sps; aux = compl\ gs\ bs$
 $ps0\ sps\ (snd\ data);$
 $remcomps = fst\ (data) \ -\ count-const-lt-components\ (fst\ aux)\ in$
 $(if\ remcomps = 0\ then$
 $full-gb\ (gs\ @\ bs)$

$hs\ data')$) \implies
 $P\ data\ bs\ ps\ (gb\text{-}schema\text{-}aux\ (rc,\ data'))\ (ab\ gs\ bs\ hs\ data')\ (ap\ gs\ bs$
 $(ps\ \text{---}\ sps)\ hs\ data')$
shows $P\ data\ bs\ ps\ (gb\text{-}schema\text{-}aux\ data\ bs\ ps)$
 $\langle proof \rangle$

lemma *gb-schema-dummy-eq-gb-schema-aux*:
assumes *struct-spec sel ap ab compl*
shows $fst\ (gb\text{-}schema\text{-}dummy\ data\ D\ bs\ ps) = gb\text{-}schema\text{-}aux\ data\ bs\ ps$
 $\langle proof \rangle$

corollary *gb-schema-aux-dgrad-p-set-le*:
assumes *dickson-grading d and struct-spec sel ap ab compl*
shows $dgrad\text{-}p\text{-}set\text{-}le\ d\ (fst\ 'set\ (gb\text{-}schema\text{-}aux\ data\ bs\ ps))\ (args\text{-}to\text{-}set\ (gs,\ bs,$
 $ps))$
 $\langle proof \rangle$

corollary *gb-schema-aux-components*:
assumes *struct-spec sel ap ab compl and set ps \subseteq set bs \times (set gs \cup set bs)*
shows $component\text{-}of\text{-}term\ 'Keys\ (fst\ 'set\ (gb\text{-}schema\text{-}aux\ data\ bs\ ps)) =$
 $component\text{-}of\text{-}term\ 'Keys\ (args\text{-}to\text{-}set\ (gs,\ bs,\ ps))$
 $\langle proof \rangle$

lemma *gb-schema-aux-pmdl*:
assumes *struct-spec sel ap ab compl and compl-pmdl compl and is-Groebner-basis*
 $(fst\ 'set\ gs)$
and $set\ ps\ \subseteq\ set\ bs\ \times\ (set\ gs\ \cup\ set\ bs)$ **and** $unique\text{-}idx\ (gs\ @\ bs)\ (snd\ data)$
and $rem\text{-}comps\text{-}spec\ (gs\ @\ bs)\ data$
shows $pmdl\ (fst\ 'set\ (gb\text{-}schema\text{-}aux\ data\ bs\ ps)) = pmdl\ (fst\ 'set\ (gs\ @\ bs))$
 $\langle proof \rangle$

corollary *gb-schema-aux-dgrad-p-set-le-init*:
assumes *dickson-grading d and struct-spec sel ap ab compl*
shows $dgrad\text{-}p\text{-}set\text{-}le\ d\ (fst\ 'set\ (gb\text{-}schema\text{-}aux\ data\ (ab\ gs\ \square\ bs\ (snd\ data))\ (ap\ gs$
 $gs\ \square\ \square\ bs\ (snd\ data))))$
 $(fst\ '(set\ gs\ \cup\ set\ bs))$
 $\langle proof \rangle$

corollary *gb-schema-aux-dgrad-p-set-init*:
assumes *dickson-grading d and struct-spec sel ap ab compl*
and $fst\ '(set\ gs\ \cup\ set\ bs) \subseteq dgrad\text{-}p\text{-}set\ d\ m$
shows $fst\ 'set\ (gb\text{-}schema\text{-}aux\ (rc,\ data)\ (ab\ gs\ \square\ bs\ data)\ (ap\ gs\ \square\ \square\ bs\ data))$
 $\subseteq dgrad\text{-}p\text{-}set\ d\ m$
 $\langle proof \rangle$

corollary *gb-schema-aux-components-init*:
assumes *struct-spec sel ap ab compl*
shows $component\text{-}of\text{-}term\ 'Keys\ (fst\ 'set\ (gb\text{-}schema\text{-}aux\ (rc,\ data)\ (ab\ gs\ \square\ bs$
 $data)\ (ap\ gs\ \square\ \square\ bs\ data))) =$

component-of-term ' Keys (fst ' set (gs @ bs))
 ⟨proof⟩

corollary *gb-schema-aux-pmdl-init:*

assumes *struct-spec sel ap ab compl and compl-pmdl compl and is-Groebner-basis (fst ' set gs)*
and *unique-idx (gs @ ab gs [] bs data) data and rem-comps-spec (gs @ ab gs [] bs data) (rc, data)*
shows *pmdl (fst ' set (gb-schema-aux (rc, data) (ab gs [] bs data) (ap gs [] [] bs data))) =*
pmdl (fst ' (set (gs @ bs)))
 ⟨proof⟩

lemma *gb-schema-aux-isGB-init:*

assumes *struct-spec sel ap ab compl and compl-conn compl and is-Groebner-basis (fst ' set gs)*
and *unique-idx (gs @ ab gs [] bs data) data and rem-comps-spec (gs @ ab gs [] bs data) (rc, data)*
shows *is-Groebner-basis (fst ' set (gb-schema-aux (rc, data) (ab gs [] bs data) (ap gs [] [] bs data)))*
 ⟨proof⟩

end

6.2.11 Functions *gb-schema-direct* and term *gb-schema-incr*

definition *gb-schema-direct* :: $(t, 'b, 'c, 'd)$ *selT* \Rightarrow $(t, 'b, 'c, 'd)$ *apT* \Rightarrow $(t, 'b, 'c, 'd)$ *abT* \Rightarrow

$(t, 'b, 'c, 'd)$ *complT* \Rightarrow $(t, 'b, 'c)$ *pdata' list* \Rightarrow $'d \Rightarrow$
 $(t, 'b::field, 'c::default)$ *pdata' list*

where *gb-schema-direct sel ap ab compl bs0 data0 =*
(let data = (length bs0, data0); bs1 = fst (add-indices (bs0, data0) (0,
data0));

bs = ab [] [] bs1 data in
map ($\lambda(f, -, d). (f, d)$)

(gb-schema-aux sel ap ab compl [] (count-rem-components bs, data)
bs (ap [] [] [] bs1 data))
)

primrec *gb-schema-incr* :: $(t, 'b, 'c, 'd)$ *selT* \Rightarrow $(t, 'b, 'c, 'd)$ *apT* \Rightarrow $(t, 'b, 'c, 'd)$ *abT* \Rightarrow

$(t, 'b, 'c, 'd)$ *complT* \Rightarrow
 $((t, 'b, 'c)$ *pdata list* \Rightarrow $(t, 'b, 'c)$ *pdata* \Rightarrow $'d \Rightarrow$ $'d) \Rightarrow$
 $(t, 'b, 'c)$ *pdata' list* \Rightarrow $'d \Rightarrow$ $(t, 'b::field, 'c::default)$

pdata' list

where

gb-schema-incr - - - - [] - = []

gb-schema-incr sel ap ab compl upd (b0 # bs) data =

(let (gs, n, data') = add-indices (gb-schema-incr sel ap ab compl upd bs data,

data) (0, *data*);
 $b = (\text{fst } b0, n, \text{snd } b0)$; $\text{data}'' = \text{upd } gs \ b \ \text{data}' \ \text{in}$
 $\text{map } (\lambda(f, -, d). (f, d))$
 $(\text{gb-schema-aux } \text{sel } \text{ap } \text{ab } \text{compl } gs \ (\text{count-rem-components } (b \# gs), \text{Suc } n, \text{data}''))$
 $(\text{ab } gs \ [] \ [b] \ (\text{Suc } n, \text{data}'')) \ (\text{ap } gs \ [] \ [] \ [b] \ (\text{Suc } n, \text{data}''))$
)

lemma (in $-$) *fst-set-drop-indices*:
 $\text{fst } ' \ (\lambda(f, -, d). (f, d)) \ ' \ A = \text{fst } ' \ A \ \text{for } A::('x \times 'y \times 'z) \ \text{set}$
 $\langle \text{proof} \rangle$

lemma *fst-gb-schema-direct*:
 $\text{fst } ' \ \text{set } (\text{gb-schema-direct } \text{sel } \text{ap } \text{ab } \text{compl } bs0 \ \text{data}0) =$
 $(\text{let } \text{data} = (\text{length } bs0, \text{data}0); \text{bs1} = \text{fst } (\text{add-indices } (bs0, \text{data}0) \ (0, \text{data}0));$
 $bs = \text{ab } [] \ [] \ \text{bs1} \ \text{data} \ \text{in}$
 $\text{fst } ' \ \text{set } (\text{gb-schema-aux } \text{sel } \text{ap } \text{ab } \text{compl } [] \ (\text{count-rem-components } bs, \text{data})$
 $bs \ (\text{ap } [] \ [] \ [] \ \text{bs1} \ \text{data}))$
)
 $\langle \text{proof} \rangle$

lemma *gb-schema-direct-dgrad-p-set*:
assumes *dickson-grading d and struct-spec sel ap ab compl and fst ' set bs \subseteq dgrad-p-set d m*
shows $\text{fst } ' \ \text{set } (\text{gb-schema-direct } \text{sel } \text{ap } \text{ab } \text{compl } bs \ \text{data}) \subseteq \text{dgrad-p-set } d \ m$
 $\langle \text{proof} \rangle$

theorem *gb-schema-direct-isGB*:
assumes *struct-spec sel ap ab compl and compl-conn compl*
shows *is-Groebner-basis (fst ' set (gb-schema-direct sel ap ab compl bs data))*
 $\langle \text{proof} \rangle$

theorem *gb-schema-direct-pmdl*:
assumes *struct-spec sel ap ab compl and compl-pmdl compl*
shows $\text{pmdl } (\text{fst } ' \ \text{set } (\text{gb-schema-direct } \text{sel } \text{ap } \text{ab } \text{compl } bs \ \text{data})) = \text{pmdl } (\text{fst } ' \ \text{set } bs)$
 $\langle \text{proof} \rangle$

lemma *fst-gb-schema-incr*:
 $\text{fst } ' \ \text{set } (\text{gb-schema-incr } \text{sel } \text{ap } \text{ab } \text{compl } \text{upd } (b0 \# bs) \ \text{data}) =$
 $(\text{let } (gs, n, \text{data}') = \text{add-indices } (\text{gb-schema-incr } \text{sel } \text{ap } \text{ab } \text{compl } \text{upd } bs \ \text{data},$
 $\text{data}) \ (0, \text{data});$
 $b = (\text{fst } b0, n, \text{snd } b0)$; $\text{data}'' = \text{upd } gs \ b \ \text{data}' \ \text{in}$
 $\text{fst } ' \ \text{set } (\text{gb-schema-aux } \text{sel } \text{ap } \text{ab } \text{compl } gs \ (\text{count-rem-components } (b \# gs),$
 $\text{Suc } n, \text{data}''))$
 $(\text{ab } gs \ [] \ [b] \ (\text{Suc } n, \text{data}'')) \ (\text{ap } gs \ [] \ [] \ [b] \ (\text{Suc } n, \text{data}''))$
)
 $\langle \text{proof} \rangle$

lemma *gb-schema-incr-dgrad-p-set*:

assumes *dickson-grading d* **and** *struct-spec sel ap ab compl*

and *fst ' set bs \subseteq dgrad-p-set d m*

shows *fst ' set (gb-schema-incr sel ap ab compl upd bs data) \subseteq dgrad-p-set d m*

<proof>

theorem *gb-schema-incr-dgrad-p-set-isGB*:

assumes *struct-spec sel ap ab compl* **and** *compl-conn compl*

shows *is-Groebner-basis (fst ' set (gb-schema-incr sel ap ab compl upd bs data))*

<proof>

theorem *gb-schema-incr-pmdl*:

assumes *struct-spec sel ap ab compl* **and** *compl-conn compl compl-pmdl compl*

shows *pmdl (fst ' set (gb-schema-incr sel ap ab compl upd bs data)) = pmdl (fst ' set bs)*

<proof>

6.3 Suitable Instances of the *add-pairs* Parameter

6.3.1 Specification of the *crit* parameters

type-synonym (in $-$) (*t, 'b, 'c, 'd*) *icritT* = $\text{nat} \times 'd \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow$

$(('t, 'b, 'c) \text{pdata list} \Rightarrow ('t, 'b, 'c) \text{pdata} \Rightarrow ('t, 'b,$

$'c) \text{pdata} \Rightarrow \text{bool}$

type-synonym (in $-$) (*t, 'b, 'c, 'd*) *ncritT* = $\text{nat} \times 'd \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow$

$(('t, 'b, 'c) \text{pdata list} \Rightarrow \text{bool} \Rightarrow$

$(\text{bool} \times ('t, 'b, 'c) \text{pdata-pair}) \text{list} \Rightarrow ('t, 'b, 'c)$

$\text{pdata} \Rightarrow$

$(('t, 'b, 'c) \text{pdata} \Rightarrow \text{bool}$

type-synonym (in $-$) (*t, 'b, 'c, 'd*) *ocritT* = $\text{nat} \times 'd \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow (\text{bool} \times ('t, 'b, 'c) \text{pdata-pair}) \text{list} \Rightarrow ('t, 'b, 'c)$

$\text{pdata} \Rightarrow$

$(('t, 'b, 'c) \text{pdata} \Rightarrow \text{bool}$

definition *icrit-spec* :: (*t, 'b::field, 'c, 'd*) *icritT* \Rightarrow *bool*

where *icrit-spec crit* \longleftrightarrow

$(\forall d m \text{ data } gs \ bs \ hs \ p \ q. \text{dickson-grading } d \longrightarrow$

$\text{fst ' (set } gs \cup \text{set } bs \cup \text{set } hs) \subseteq \text{dgrad-p-set } d \ m \longrightarrow \text{unique-idx } (gs \ @$

$bs \ @ \ hs) \ \text{data} \longrightarrow$

$\text{is-Groebner-basis } (\text{fst ' set } gs) \longrightarrow p \in \text{set } hs \longrightarrow q \in \text{set } gs \cup \text{set } bs \cup$

$\text{set } hs \longrightarrow$

$\text{fst } p \neq 0 \longrightarrow \text{fst } q \neq 0 \longrightarrow \text{crit } \text{data } gs \ bs \ hs \ p \ q \longrightarrow$

$\text{crit-pair-cbelow-on } d \ m \ (\text{fst ' (set } gs \cup \text{set } bs \cup \text{set } hs)) \ (\text{fst } p) \ (\text{fst } q))$

Criteria satisfying *icrit-spec* can be used for discarding pairs *instantly*, without reference to any other pairs. The product criterion for scalar polyno-

mials satisfies *icrit-spec*, and so does the component criterion (which checks whether the component-indices of the leading terms of two polynomials are identical).

definition *ncrit-spec* :: ('t, 'b::field, 'c, 'd) *ncritT* ⇒ bool

where *ncrit-spec crit* ⇔
 $(\forall d\ m\ data\ gs\ bs\ hs\ ps\ B\ q\text{-in-}bs\ p\ q.\ d\text{ickson-grad}\ d \longrightarrow set\ gs \cup set\ bs \cup set\ hs \subseteq B \longrightarrow$
 $fst\ 'B \subseteq d\text{grad-p-set}\ d\ m \longrightarrow snd\ 'set\ ps \subseteq set\ hs \times (set\ gs \cup set\ bs \cup set\ hs) \longrightarrow$
 $unique\text{-idx}\ (gs\ @\ bs\ @\ hs)\ data \longrightarrow is\text{-Groebner-basis}\ (fst\ 'set\ gs) \longrightarrow$
 $(q\ \text{in}\ bs \longrightarrow (q \in set\ gs \cup set\ bs)) \longrightarrow$
 $(\forall p'\ q'. (p', q') \in_p\ snd\ 'set\ ps \longrightarrow fst\ p' \neq 0 \longrightarrow fst\ q' \neq 0 \longrightarrow$
 $crit\text{-pair-cbelow-on}\ d\ m\ (fst\ 'B)\ (fst\ p')\ (fst\ q')) \longrightarrow$
 $(\forall p'\ q'. p' \in set\ gs \cup set\ bs \longrightarrow q' \in set\ gs \cup set\ bs \longrightarrow fst\ p' \neq 0 \longrightarrow$
 $fst\ q' \neq 0 \longrightarrow$
 $crit\text{-pair-cbelow-on}\ d\ m\ (fst\ 'B)\ (fst\ p')\ (fst\ q')) \longrightarrow$
 $p \in set\ hs \longrightarrow q \in set\ gs \cup set\ bs \cup set\ hs \longrightarrow fst\ p \neq 0 \longrightarrow fst\ q \neq 0$
 \longrightarrow
 $crit\ data\ gs\ bs\ hs\ q\text{-in-}bs\ ps\ p\ q \longrightarrow$
 $crit\text{-pair-cbelow-on}\ d\ m\ (fst\ 'B)\ (fst\ p)\ (fst\ q))$

definition *ocrit-spec* :: ('t, 'b::field, 'c, 'd) *ocritT* ⇒ bool

where *ocrit-spec crit* ⇔
 $(\forall d\ m\ data\ hs\ ps\ B\ p\ q.\ d\text{ickson-grad}\ d \longrightarrow set\ hs \subseteq B \longrightarrow fst\ 'B \subseteq$
 $d\text{grad-p-set}\ d\ m \longrightarrow$
 $unique\text{-idx}\ (p\ \#\ q\ \#\ hs\ @\ (map\ (fst\ o\ snd)\ ps)\ @\ (map\ (snd\ o\ snd)\ ps))\ data \longrightarrow$
 $(\forall p'\ q'. (p', q') \in_p\ snd\ 'set\ ps \longrightarrow fst\ p' \neq 0 \longrightarrow fst\ q' \neq 0 \longrightarrow$
 $crit\text{-pair-cbelow-on}\ d\ m\ (fst\ 'B)\ (fst\ p')\ (fst\ q')) \longrightarrow$
 $p \in B \longrightarrow q \in B \longrightarrow fst\ p \neq 0 \longrightarrow fst\ q \neq 0 \longrightarrow$
 $crit\ data\ hs\ ps\ p\ q \longrightarrow crit\text{-pair-cbelow-on}\ d\ m\ (fst\ 'B)\ (fst\ p)\ (fst\ q))$

Criteria satisfying *ncrit-spec* can be used for discarding new pairs by reference to new and old elements, whereas criteria satisfying *ocrit-spec* can be used for discarding old pairs by reference to new elements *only* (no existing ones!). The chain criterion satisfies both *ncrit-spec* and *ocrit-spec*.

lemma *icrit-specI*:

assumes $\bigwedge d\ m\ data\ gs\ bs\ hs\ p\ q.$
 $d\text{ickson-grad}\ d \Longrightarrow fst\ '(set\ gs \cup set\ bs \cup set\ hs) \subseteq d\text{grad-p-set}\ d\ m$
 \Longrightarrow
 $unique\text{-idx}\ (gs\ @\ bs\ @\ hs)\ data \Longrightarrow is\text{-Groebner-basis}\ (fst\ 'set\ gs) \Longrightarrow$
 $p \in set\ hs \Longrightarrow q \in set\ gs \cup set\ bs \cup set\ hs \Longrightarrow fst\ p \neq 0 \Longrightarrow fst\ q \neq 0$
 \Longrightarrow
 $crit\ data\ gs\ bs\ hs\ p\ q \Longrightarrow$
 $crit\text{-pair-cbelow-on}\ d\ m\ (fst\ '(set\ gs \cup set\ bs \cup set\ hs))\ (fst\ p)\ (fst\ q)$
shows *icrit-spec crit*
 ⟨proof⟩

lemma *icrit-specD*:

assumes *icrit-spec crit and dickson-grading d*
and $\text{fst } \text{' } (set\ gs \cup set\ bs \cup set\ hs) \subseteq dgrad\text{-}p\text{-}set\ d\ m$ **and** *unique-idx (gs @ bs @ hs) data*
and *is-Groebner-basis (fst ' set gs) and $p \in set\ hs$ and $q \in set\ gs \cup set\ bs \cup set\ hs$*
and $\text{fst } p \neq 0$ **and** $\text{fst } q \neq 0$ **and** *crit data gs bs hs p q*
shows *crit-pair-cbelow-on d m (fst ' (set gs \cup set bs \cup set hs)) (fst p) (fst q)*
<proof>

lemma *ncrit-specI*:

assumes $\bigwedge d\ m\ data\ gs\ bs\ hs\ ps\ B\ q\text{-in-}bs\ p\ q.$
dickson-grading d $\implies set\ gs \cup set\ bs \cup set\ hs \subseteq B \implies$
 $\text{fst } \text{' } B \subseteq dgrad\text{-}p\text{-}set\ d\ m \implies \text{snd } \text{' } set\ ps \subseteq set\ hs \times (set\ gs \cup set\ bs$
 $\cup set\ hs) \implies$
unique-idx (gs @ bs @ hs) data $\implies is\text{-}Groebner\text{-}basis\ (fst\ \text{' } set\ gs) \implies$
(q-in-bs $\longrightarrow q \in set\ gs \cup set\ bs) \implies$
($\bigwedge p' q'. (p', q') \in_p \text{snd } \text{' } set\ ps \implies \text{fst } p' \neq 0 \implies \text{fst } q' \neq 0 \implies$
crit-pair-cbelow-on d m (fst ' B) (fst p') (fst q')) \implies
($\bigwedge p' q'. p' \in set\ gs \cup set\ bs \implies q' \in set\ gs \cup set\ bs \implies \text{fst } p' \neq 0 \implies$
 $\text{fst } q' \neq 0 \implies$
crit-pair-cbelow-on d m (fst ' B) (fst p') (fst q')) \implies
 $p \in set\ hs \implies q \in set\ gs \cup set\ bs \cup set\ hs \implies \text{fst } p \neq 0 \implies \text{fst } q \neq 0$
 \implies
crit data gs bs hs q-in-bs ps p q \implies
crit-pair-cbelow-on d m (fst ' B) (fst p) (fst q)
shows *ncrit-spec crit*
<proof>

lemma *ncrit-specD*:

assumes *ncrit-spec crit and dickson-grading d and $set\ gs \cup set\ bs \cup set\ hs \subseteq B$*
and $\text{fst } \text{' } B \subseteq dgrad\text{-}p\text{-}set\ d\ m$ **and** $\text{snd } \text{' } set\ ps \subseteq set\ hs \times (set\ gs \cup set\ bs \cup set\ hs)$
and *unique-idx (gs @ bs @ hs) data and is-Groebner-basis (fst ' set gs)*
and *q-in-bs $\implies q \in set\ gs \cup set\ bs$*
and $\bigwedge p' q'. (p', q') \in_p \text{snd } \text{' } set\ ps \implies \text{fst } p' \neq 0 \implies \text{fst } q' \neq 0 \implies$
crit-pair-cbelow-on d m (fst ' B) (fst p') (fst q')
and $\bigwedge p' q'. p' \in set\ gs \cup set\ bs \implies q' \in set\ gs \cup set\ bs \implies \text{fst } p' \neq 0 \implies \text{fst } q' \neq 0 \implies$
crit-pair-cbelow-on d m (fst ' B) (fst p') (fst q')
and $p \in set\ hs$ **and** $q \in set\ gs \cup set\ bs \cup set\ hs$ **and** $\text{fst } p \neq 0$ **and** $\text{fst } q \neq 0$
and *crit data gs bs hs q-in-bs ps p q*
shows *crit-pair-cbelow-on d m (fst ' B) (fst p) (fst q)*
<proof>

lemma *ocrit-specI*:

assumes $\bigwedge d\ m\ data\ hs\ ps\ B\ p\ q.$
dickson-grading d $\implies set\ hs \subseteq B \implies \text{fst } \text{' } B \subseteq dgrad\text{-}p\text{-}set\ d\ m \implies$
unique-idx (p # q # hs @ (map (fst \circ snd) ps) @ (map (snd \circ snd)

$ps)) \text{ data} \implies$
 $(\bigwedge p' q'. (p', q') \in_p \text{snd} \text{ ' set } ps \implies \text{fst } p' \neq 0 \implies \text{fst } q' \neq 0 \implies$
 $\text{crit-pair-cbelow-on } d \ m \ (\text{fst} \text{ ' } B) \ (\text{fst } p') \ (\text{fst } q')) \implies$
 $p \in B \implies q \in B \implies \text{fst } p \neq 0 \implies \text{fst } q \neq 0 \implies$
 $\text{crit data } hs \ ps \ p \ q \implies \text{crit-pair-cbelow-on } d \ m \ (\text{fst} \text{ ' } B) \ (\text{fst } p) \ (\text{fst } q)$
shows *ocrit-spec crit*
 $\langle \text{proof} \rangle$

lemma *ocrit-specD*:

assumes *ocrit-spec crit and dickson-grading d and set* $hs \subseteq B$ **and** $\text{fst} \text{ ' } B \subseteq$
 $d\text{grad-p-set } d \ m$
and *unique-idx* $(p \# q \# hs \ @ \ (\text{map} \ (\text{fst} \circ \text{snd}) \ ps) \ @ \ (\text{map} \ (\text{snd} \circ \text{snd}) \ ps))$
 data
and $\bigwedge p' q'. (p', q') \in_p \text{snd} \text{ ' set } ps \implies \text{fst } p' \neq 0 \implies \text{fst } q' \neq 0 \implies$
 $\text{crit-pair-cbelow-on } d \ m \ (\text{fst} \text{ ' } B) \ (\text{fst } p') \ (\text{fst } q')$
and $p \in B$ **and** $q \in B$ **and** $\text{fst } p \neq 0$ **and** $\text{fst } q \neq 0$
and *crit data* $hs \ ps \ p \ q$
shows *crit-pair-cbelow-on* $d \ m \ (\text{fst} \text{ ' } B) \ (\text{fst } p) \ (\text{fst } q)$
 $\langle \text{proof} \rangle$

6.3.2 Suitable instances of the *crit* parameters

definition *component-crit* :: $(t, 'b::\text{zero}, 'c, 'd)$ *icritT*

where *component-crit data* $gs \ bs \ hs \ p \ q \longleftrightarrow (\text{component-of-term} \ (\text{lt} \ (\text{fst } p)) \neq$
 $\text{component-of-term} \ (\text{lt} \ (\text{fst } q)))$

lemma *icrit-spec-component-crit*: *icrit-spec* $(\text{component-crit}::(t, 'b::\text{field}, 'c, 'd)$
 $\text{icritT})$
 $\langle \text{proof} \rangle$

The product criterion is only applicable to scalar polynomials.

definition *product-crit* :: $(a, 'b::\text{zero}, 'c, 'd)$ *icritT*

where *product-crit data* $gs \ bs \ hs \ p \ q \longleftrightarrow (\text{gcd} \ (\text{punit.lt} \ (\text{fst } p)) \ (\text{punit.lt} \ (\text{fst } q))$
 $= 0)$

lemma **(in** *gd-term*) *icrit-spec-product-crit*: *punit.icrit-spec* $(\text{product-crit}::(a, 'b::\text{field},$
 $'c, 'd)$ *icritT*)
 $\langle \text{proof} \rangle$

component-crit and *product-crit* ignore the *data* parameter.

fun **(in** $-)$ *pair-in-list* :: $(\text{bool} \times (a, 'b, 'c)$ *pdata-pair*) *list* $\Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{bool}$
where

$\text{pair-in-list} \ [] \ - \ = \ \text{False}$
 $\text{pair-in-list} \ ((-, (-, i', -), (-, j', -)) \ # \ ps) \ i \ j \ =$
 $((i = i' \wedge j = j') \vee (i = j' \wedge j = i') \vee \text{pair-in-list } ps \ i \ j)$

lemma **(in** $-)$ *pair-in-listE*:

assumes *pair-in-list* $ps \ i \ j$

obtains $p \ q \ a \ b$ **where** $((p, i, a), (q, j, b)) \in_p \text{snd} \text{ ' set } ps$

$\langle \text{proof} \rangle$

definition $\text{chain-ncrit} :: ('t, 'b::\text{zero}, 'c, 'd) \text{ncrit}T$

where $\text{chain-ncrit data gs bs hs q-in-bs ps p q} \longleftrightarrow$
 $(\text{let } v = \text{lt } (\text{fst } p); l = \text{term-of-pair } (\text{lcs } (\text{pp-of-term } v) (\text{lp } (\text{fst } q))),$
 $\text{component-of-term } v);$
 $i = \text{fst } (\text{snd } p); j = \text{fst } (\text{snd } q) \text{ in}$
 $(\exists r \in \text{set } gs. \text{let } k = \text{fst } (\text{snd } r) \text{ in}$
 $k \neq i \wedge k \neq j \wedge \text{lt } (\text{fst } r) \text{ adds}_t l \wedge \text{pair-in-list } ps \ i \ k \wedge (q\text{-in-bs} \vee$
 $\text{pair-in-list } ps \ j \ k) \wedge \text{fst } r \neq 0) \vee$
 $(\exists r \in \text{set } bs. \text{let } k = \text{fst } (\text{snd } r) \text{ in}$
 $k \neq i \wedge k \neq j \wedge \text{lt } (\text{fst } r) \text{ adds}_t l \wedge \text{pair-in-list } ps \ i \ k \wedge (q\text{-in-bs} \vee$
 $\text{pair-in-list } ps \ j \ k) \wedge \text{fst } r \neq 0) \vee$
 $(\exists h \in \text{set } hs. \text{let } k = \text{fst } (\text{snd } h) \text{ in}$
 $k \neq i \wedge k \neq j \wedge \text{lt } (\text{fst } h) \text{ adds}_t l \wedge \text{pair-in-list } ps \ i \ k \wedge \text{pair-in-list}$
 $ps \ j \ k \wedge \text{fst } h \neq 0))$

definition $\text{chain-ocrit} :: ('t, 'b::\text{zero}, 'c, 'd) \text{ocrit}T$

where $\text{chain-ocrit data hs ps p q} \longleftrightarrow$
 $(\text{let } v = \text{lt } (\text{fst } p); l = \text{term-of-pair } (\text{lcs } (\text{pp-of-term } v) (\text{lp } (\text{fst } q))),$
 $\text{component-of-term } v);$
 $i = \text{fst } (\text{snd } p); j = \text{fst } (\text{snd } q) \text{ in}$
 $(\exists h \in \text{set } hs. \text{let } k = \text{fst } (\text{snd } h) \text{ in}$
 $k \neq i \wedge k \neq j \wedge \text{lt } (\text{fst } h) \text{ adds}_t l \wedge \text{pair-in-list } ps \ i \ k \wedge \text{pair-in-list}$
 $ps \ j \ k \wedge \text{fst } h \neq 0))$

chain-ncrit and chain-ocrit ignore the data parameter.

lemma chain-ncritE :

assumes $\text{chain-ncrit data gs bs hs q-in-bs ps p q}$ **and** $\text{snd } ' \text{ set } ps \subseteq \text{set } hs \times$
 $(\text{set } gs \cup \text{set } bs \cup \text{set } hs)$
and $\text{unique-idx } (gs \ @ \ bs \ @ \ hs) \ \text{data}$ **and** $p \in \text{set } hs$ **and** $q \in \text{set } gs \cup \text{set } bs \cup$
 $\text{set } hs$
obtains r **where** $r \in \text{set } gs \cup \text{set } bs \cup \text{set } hs$ **and** $\text{fst } r \neq 0$ **and** $r \neq p$ **and** r
 $\neq q$
and $\text{lt } (\text{fst } r) \text{ adds}_t \text{term-of-pair } (\text{lcs } (\text{lp } (\text{fst } p)) (\text{lp } (\text{fst } q))), \text{component-of-term}$
 $(\text{lt } (\text{fst } p))$
and $(p, r) \in_p \text{snd } ' \text{ set } ps$ **and** $(r \in \text{set } gs \cup \text{set } bs \wedge q\text{-in-bs}) \vee (q, r) \in_p \text{snd}$
 $' \text{ set } ps$
 $\langle \text{proof} \rangle$

lemma chain-ocritE :

assumes $\text{chain-ocrit data hs ps p q}$
and $\text{unique-idx } (p \ # \ q \ # \ hs \ @ \ (\text{map } (\text{fst } \circ \ \text{snd}) \ ps) \ @ \ (\text{map } (\text{snd } \circ \ \text{snd}) \ ps))$
 data **(is** $\text{unique-idx } ?xs \ -)$
obtains h **where** $h \in \text{set } hs$ **and** $\text{fst } h \neq 0$ **and** $h \neq p$ **and** $h \neq q$
and $\text{lt } (\text{fst } h) \text{ adds}_t \text{term-of-pair } (\text{lcs } (\text{lp } (\text{fst } p)) (\text{lp } (\text{fst } q))), \text{component-of-term}$
 $(\text{lt } (\text{fst } p))$
and $(p, h) \in_p \text{snd } ' \text{ set } ps$ **and** $(q, h) \in_p \text{snd } ' \text{ set } ps$
 $\langle \text{proof} \rangle$

lemma *ncrit-spec-chain-ncrit*: *ncrit-spec* (*chain-ncrit*::('t, 'b::field, 'c, 'd) *ncritT*)
 ⟨*proof*⟩

lemma *ocrit-spec-chain-ocrit*: *ocrit-spec* (*chain-ocrit*::('t, 'b::field, 'c, 'd) *ocritT*)
 ⟨*proof*⟩

lemma *icrit-spec-no-crit*: *icrit-spec* ((λ- - - - - . *False*)::('t, 'b::field, 'c, 'd) *icritT*)
 ⟨*proof*⟩

lemma *ncrit-spec-no-crit*: *ncrit-spec* ((λ- - - - - . *False*)::('t, 'b::field, 'c, 'd) *ncritT*)
 ⟨*proof*⟩

lemma *ocrit-spec-no-crit*: *ocrit-spec* ((λ- - - - - . *False*)::('t, 'b::field, 'c, 'd) *ocritT*)
 ⟨*proof*⟩

6.3.3 Creating Initial List of New Pairs

type-synonym (**in** -) ('t, 'b, 'c) *apsT* = *bool* ⇒ ('t, 'b, 'c) *pdata list* ⇒ ('t, 'b, 'c) *pdata list* ⇒
 ('t, 'b, 'c) *pdata* ⇒ (*bool* × ('t, 'b, 'c) *pdata-pair*) *list*
 ⇒
 (*bool* × ('t, 'b, 'c) *pdata-pair*) *list*

type-synonym (**in** -) ('t, 'b, 'c, 'd) *npT* = ('t, 'b, 'c) *pdata list* ⇒ ('t, 'b, 'c) *pdata list* ⇒
 ('t, 'b, 'c) *pdata list* ⇒ *nat* × 'd ⇒
 (*bool* × ('t, 'b, 'c) *pdata-pair*) *list*

definition *np-spec* :: ('t, 'b, 'c, 'd) *npT* ⇒ *bool*
where *np-spec np* ⇔ (∀ *gs bs hs data*.
snd ' *set* (*np gs bs hs data*) ⊆ *set hs* × (*set gs* ∪ *set bs* ∪ *set hs*) ∧
set hs × (*set gs* ∪ *set bs*) ⊆ *snd* ' *set* (*np gs bs hs data*) ∧
 (∀ *a b*. *a* ∈ *set hs* → *b* ∈ *set hs* → *a* ≠ *b* → (*a*, *b*) ∈_{*p*} *snd* ' *set* (*np gs bs hs data*)) ∧
 (∀ *p q*. (*True*, *p*, *q*) ∈ *set* (*np gs bs hs data*) → *q* ∈ *set gs* ∪ *set bs*))

lemma *np-specI*:
assumes ∧*gs bs hs data*.
snd ' *set* (*np gs bs hs data*) ⊆ *set hs* × (*set gs* ∪ *set bs* ∪ *set hs*) ∧
set hs × (*set gs* ∪ *set bs*) ⊆ *snd* ' *set* (*np gs bs hs data*) ∧
 (∀ *a b*. *a* ∈ *set hs* → *b* ∈ *set hs* → *a* ≠ *b* → (*a*, *b*) ∈_{*p*} *snd* ' *set* (*np gs bs hs data*)) ∧
 (∀ *p q*. (*True*, *p*, *q*) ∈ *set* (*np gs bs hs data*) → *q* ∈ *set gs* ∪ *set bs*)
shows *np-spec np*
 ⟨*proof*⟩

lemma *np-specD1*:

assumes *np-spec np*

shows $\text{snd} \text{ ' set } (np \text{ gs } bs \text{ hs } data) \subseteq \text{set } hs \times (\text{set } gs \cup \text{set } bs \cup \text{set } hs)$

<proof>

lemma *np-specD2*:

assumes *np-spec np*

shows $\text{set } hs \times (\text{set } gs \cup \text{set } bs) \subseteq \text{snd} \text{ ' set } (np \text{ gs } bs \text{ hs } data)$

<proof>

lemma *np-specD3*:

assumes *np-spec np* **and** $a \in \text{set } hs$ **and** $b \in \text{set } hs$ **and** $a \neq b$

shows $(a, b) \in_p \text{snd} \text{ ' set } (np \text{ gs } bs \text{ hs } data)$

<proof>

lemma *np-specD4*:

assumes *np-spec np* **and** $(\text{True}, p, q) \in \text{set } (np \text{ gs } bs \text{ hs } data)$

shows $q \in \text{set } gs \cup \text{set } bs$

<proof>

lemma *np-specE*:

assumes *np-spec np* **and** $p \in \text{set } hs$ **and** $q \in \text{set } gs \cup \text{set } bs \cup \text{set } hs$ **and** $p \neq q$

assumes 1: $\bigwedge q \text{-in-} bs. (q \text{-in-} bs, p, q) \in \text{set } (np \text{ gs } bs \text{ hs } data) \implies \text{thesis}$

assumes 2: $\bigwedge p \text{-in-} bs. (p \text{-in-} bs, q, p) \in \text{set } (np \text{ gs } bs \text{ hs } data) \implies \text{thesis}$

shows *thesis*

<proof>

definition *add-pairs-single-naive* :: $'d \Rightarrow ('t, 'b::\text{zero}, 'c) \text{ apsT}$

where *add-pairs-single-naive data flag gs bs h ps* = $ps \text{ @ } (\text{map } (\lambda g. (\text{flag}, h, g)) \text{ gs}) \text{ @ } (\text{map } (\lambda b. (\text{flag}, h, b)) \text{ bs})$

lemma *set-add-pairs-single-naive*:

$\text{set } (\text{add-pairs-single-naive data flag gs bs h ps}) = \text{set } ps \cup \text{Pair flag ' } (\{h\} \times (\text{set } gs \cup \text{set } bs))$

<proof>

fun *add-pairs-single-sorted* :: $((\text{bool} \times ('t, 'b, 'c) \text{ pdata-pair}) \Rightarrow (\text{bool} \times ('t, 'b, 'c) \text{ pdata-pair})) \Rightarrow \text{bool} \Rightarrow$

$(('t, 'b::\text{zero}, 'c) \text{ apsT} \text{ where}$

add-pairs-single-sorted - - $\square \square - ps = ps \mid$

add-pairs-single-sorted rel flag $\square (b \# bs) h ps =$

add-pairs-single-sorted rel flag $\square bs h (\text{insort-wrt rel } (\text{flag}, h, b) ps) \mid$

add-pairs-single-sorted rel flag $(g \# gs) bs h ps =$

add-pairs-single-sorted rel flag gs bs h (\text{insort-wrt rel } (\text{flag}, h, g) ps)

lemma *set-add-pairs-single-sorted*:

$\text{set } (\text{add-pairs-single-sorted rel flag gs bs h ps}) = \text{set } ps \cup \text{Pair flag ' } (\{h\} \times (\text{set } gs \cup \text{set } bs))$

<proof>

primrec (**in** $-$) *pairs* :: ('t, 'b, 'c) *apsT* \Rightarrow *bool* \Rightarrow ('t, 'b, 'c) *pdata list* \Rightarrow (*bool* \times ('t, 'b, 'c) *pdata-pair*) *list*

where

pairs - - [] = []

pairs *aps* *flag* ($x \# xs$) = *aps* *flag* [] *xs* *x* (*pairs* *aps* *flag* *xs*)

lemma *pairs-subset*:

assumes $\bigwedge gs\ bs\ h\ ps.$ *set* (*aps* *flag* *gs* *bs* *h* *ps*) = *set* *ps* \cup *Pair* *flag* ' ($\{h\} \times$ (*set* *gs* \cup *set* *bs*))

shows *set* (*pairs* *aps* *flag* *xs*) \subseteq *Pair* *flag* ' (*set* *xs* \times *set* *xs*)

<proof>

lemma *in-pairsI*:

assumes $\bigwedge gs\ bs\ h\ ps.$ *set* (*aps* *flag* *gs* *bs* *h* *ps*) = *set* *ps* \cup *Pair* *flag* ' ($\{h\} \times$ (*set* *gs* \cup *set* *bs*))

and $a \neq b$ **and** $a \in$ *set* *xs* **and** $b \in$ *set* *xs*

shows (*flag*, *a*, *b*) \in *set* (*pairs* *aps* *flag* *xs*) \vee (*flag*, *b*, *a*) \in *set* (*pairs* *aps* *flag* *xs*)

<proof>

corollary *in-pairsI'*:

assumes $\bigwedge gs\ bs\ h\ ps.$ *set* (*aps* *flag* *gs* *bs* *h* *ps*) = *set* *ps* \cup *Pair* *flag* ' ($\{h\} \times$ (*set* *gs* \cup *set* *bs*))

and $a \in$ *set* *xs* **and** $b \in$ *set* *xs* **and** $a \neq b$

shows (*a*, *b*) \in_p *snd* ' *set* (*pairs* *aps* *flag* *xs*)

<proof>

definition *new-pairs-naive* :: ('t, 'b::zero, 'c, 'd) *npT*

where *new-pairs-naive* *gs* *bs* *hs* *data* =

fold (*add-pairs-single-naive* *data* *True* *gs* *bs*) *hs* (*pairs* (*add-pairs-single-naive* *data*) *False* *hs*)

definition *new-pairs-sorted* :: (*nat* \times 'd \Rightarrow (*bool* \times ('t, 'b, 'c) *pdata-pair*) \Rightarrow (*bool* \times ('t, 'b, 'c) *pdata-pair*) \Rightarrow *bool*) \Rightarrow

('t, 'b::zero, 'c, 'd) *npT*

where *new-pairs-sorted* *rel* *gs* *bs* *hs* *data* =

fold (*add-pairs-single-sorted* (*rel* *data*) *True* *gs* *bs*) *hs* (*pairs* (*add-pairs-single-sorted* (*rel* *data*)) *False* *hs*)

lemma *set-fold-aps*:

assumes $\bigwedge gs\ bs\ h\ ps.$ *set* (*aps* *flag* *gs* *bs* *h* *ps*) = *set* *ps* \cup *Pair* *flag* ' ($\{h\} \times$ (*set* *gs* \cup *set* *bs*))

shows *set* (*fold* (*aps* *flag* *gs* *bs*) *hs* *ps*) = *Pair* *flag* ' (*set* *hs* \times (*set* *gs* \cup *set* *bs*)) \cup *set* *ps*

<proof>

lemma *set-new-pairs-naive*:

set (*new-pairs-naive* *gs* *bs* *hs* *data*) =

Pair True ' (set hs × (set gs ∪ set bs)) ∪ set (pairs (add-pairs-single-naive data) False hs)
 ⟨proof⟩

lemma *set-new-pairs-sorted*:

set (new-pairs-sorted rel gs bs hs data) =
Pair True ' (set hs × (set gs ∪ set bs)) ∪ set (pairs (add-pairs-single-sorted (rel data)) False hs)
 ⟨proof⟩

lemma (*in -*) *fst-snd-Pair [simp]*:

shows *fst ∘ Pair x = (λ-. x)* **and** *snd ∘ Pair x = id*
 ⟨proof⟩

lemma *np-spec-new-pairs-naive*: *np-spec new-pairs-naive*

⟨proof⟩

lemma *np-spec-new-pairs-sorted*: *np-spec (new-pairs-sorted rel)*

⟨proof⟩

new-pairs-naive gs bs hs data and *new-pairs-sorted rel gs bs hs data* return lists of triples (*q-in-bs*, *p*, *q*), where *q-in-bs* indicates whether *q* is contained in the list *gs @ bs* or in the list *hs*. *p* is always contained in *hs*.

definition *canon-pair-order-aux* :: (*t*, *'b*::*zero*, *'c*) *pdata-pair* ⇒ (*t*, *'b*, *'c*) *pdata-pair* ⇒ *bool*

where *canon-pair-order-aux p q* ⇔
 (*lcs (lp (fst (fst p))) (lp (fst (snd p)))* ≤ *lcs (lp (fst (fst q))) (lp (fst (snd q)))*)

abbreviation *canon-pair-order data p q* ≡ *canon-pair-order-aux (snd p) (snd q)*

abbreviation *canon-pair-comb* ≡ *merge-wrt canon-pair-order-aux*

6.3.4 Applying Criteria to New Pairs

definition *apply-icrit* :: (*t*, *'b*, *'c*, *'d*) *icritT* ⇒ (*nat* × *'d*) ⇒ (*t*, *'b*, *'c*) *pdata list* ⇒

(*t*, *'b*, *'c*) *pdata list* ⇒ (*t*, *'b*, *'c*) *pdata list* ⇒
 (*bool* × (*t*, *'b*, *'c*) *pdata-pair*) *list* ⇒
 (*bool* × *bool* × (*t*, *'b*, *'c*) *pdata-pair*) *list*

where *apply-icrit crit data gs bs hs ps* = (*let c = crit data gs bs hs in map*
 (*λ(q-in-bs, p, q). (c p q, q-in-bs, p, q)*) *ps*)

lemma *fst-apply-icrit*:

assumes *icrit-spec crit* **and** *dickson-grading d*

and *fst ' (set gs ∪ set bs ∪ set hs) ⊆ dgrad-p-set d m* **and** *unique-idx (gs @ bs @ hs) data*

and *is-Groebner-basis (fst ' set gs)* **and** *p ∈ set hs* **and** *q ∈ set gs ∪ set bs ∪ set hs*

and $\text{fst } p \neq 0$ **and** $\text{fst } q \neq 0$ **and** $(\text{True}, q\text{-in-}bs, p, q) \in \text{set } (\text{apply-icrit crit data } gs \ bs \ hs \ ps)$

shows $\text{crit-pair-cbelow-on } d \ m \ (\text{fst } \cdot \ (\text{set } gs \cup \text{set } bs \cup \text{set } hs)) \ (\text{fst } p) \ (\text{fst } q)$
 $\langle \text{proof} \rangle$

lemma $\text{snd-apply-icrit [simp]}$: $\text{map snd } (\text{apply-icrit crit data } gs \ bs \ hs \ ps) = ps$
 $\langle \text{proof} \rangle$

lemma $\text{set-snd-apply-icrit [simp]}$: $\text{snd } \cdot \ \text{set } (\text{apply-icrit crit data } gs \ bs \ hs \ ps) = \text{set } ps$
 $\langle \text{proof} \rangle$

definition $\text{apply-ncrit} :: ('t, 'b, 'c, 'd) \text{ncrit}T \Rightarrow (\text{nat} \times 'd) \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow$

$(('t, 'b, 'c) \text{pdata list} \Rightarrow ('t, 'b, 'c) \text{pdata list} \Rightarrow$
 $(\text{bool} \times \text{bool} \times ('t, 'b, 'c) \text{pdata-pair}) \text{list} \Rightarrow$
 $(\text{bool} \times ('t, 'b, 'c) \text{pdata-pair}) \text{list}$

where $\text{apply-ncrit crit data } gs \ bs \ hs \ ps =$

$(\text{let } c = \text{crit data } gs \ bs \ hs \ \text{in}$

$\text{rev } (\text{fold } (\lambda(ic, q\text{-in-}bs, p, q). \lambda ps'. \text{if } \neg ic \wedge c \ q\text{-in-}bs \ ps' \ p \ q \ \text{then } ps'$

$\text{else } (ic, p, q) \# ps') \ ps \ \square))$

lemma $\text{apply-ncrit-append}$:

$\text{apply-ncrit crit data } gs \ bs \ hs \ (xs \ @ \ ys) =$

$\text{rev } (\text{fold } (\lambda(ic, q\text{-in-}bs, p, q). \lambda ps'. \text{if } \neg ic \wedge \text{crit data } gs \ bs \ hs \ q\text{-in-}bs \ ps' \ p \ q$

$\text{then } ps' \ \text{else } (ic, p, q) \# ps') \ ys$

$(\text{rev } (\text{apply-ncrit crit data } gs \ bs \ hs \ xs)))$

$\langle \text{proof} \rangle$

lemma fold-superset :

$\text{set } acc \subseteq$

$\text{set } (\text{fold } (\lambda(ic, q\text{-in-}bs, p, q). \lambda ps'. \text{if } \neg ic \wedge c \ q\text{-in-}bs \ ps' \ p \ q \ \text{then } ps' \ \text{else } (ic,$

$p, q) \# ps') \ ps \ acc)$

$\langle \text{proof} \rangle$

lemma $\text{apply-ncrit-superset}$:

$\text{set } (\text{apply-ncrit crit data } gs \ bs \ hs \ ps) \subseteq \text{set } (\text{apply-ncrit crit data } gs \ bs \ hs \ (ps \ @$

$qs)) \ (\text{is } ?l \subseteq ?r)$

$\langle \text{proof} \rangle$

lemma $\text{apply-ncrit-subset-aux}$:

assumes $(ic, p, q) \in \text{set } (\text{fold}$

$(\lambda(ic, q\text{-in-}bs, p, q). \lambda ps'. \text{if } \neg ic \wedge c \ q\text{-in-}bs \ ps' \ p \ q \ \text{then } ps' \ \text{else } (ic, p,$

$q) \# ps') \ ps \ acc)$

shows $(ic, p, q) \in \text{set } acc \vee (\exists q\text{-in-}bs. (ic, q\text{-in-}bs, p, q) \in \text{set } ps)$

$\langle \text{proof} \rangle$

corollary $\text{apply-ncrit-subset}$:

assumes $(ic, p, q) \in \text{set } (\text{apply-ncrit crit data } gs \ bs \ hs \ ps)$

obtains $q\text{-in-}bs$ **where** $(ic, q\text{-in-}bs, p, q) \in set\ ps$
 $\langle proof \rangle$

corollary $apply\text{-ncrit}\text{-subset}'$: $snd\ ' set\ (apply\text{-ncrit}\ crit\ data\ gs\ bs\ hs\ ps) \subseteq snd\ ' set\ ps$
 $\langle proof \rangle$

lemma $not\text{-in}\text{-apply}\text{-ncrit}$:

assumes $(ic, p, q) \notin set\ (apply\text{-ncrit}\ crit\ data\ gs\ bs\ hs\ (xs\ @\ ((ic, q\text{-in-}bs, p, q)\ #\ ys)))$
shows $crit\ data\ gs\ bs\ hs\ q\text{-in-}bs\ (rev\ (apply\text{-ncrit}\ crit\ data\ gs\ bs\ hs\ xs))\ p\ q$
 $\langle proof \rangle$

lemma $(in\ -)\ setE$:

assumes $x \in set\ xs$
obtains $ys\ zs$ **where** $xs = ys\ @\ (x\ \#\ zs)$
 $\langle proof \rangle$

lemma $apply\text{-ncrit}\text{-connectible}$:

assumes $ncrit\text{-spec}\ crit$ **and** $dickson\text{-grading}\ d$
and $set\ gs \cup set\ bs \cup set\ hs \subseteq B$ **and** $fst\ ' B \subseteq dgrad\text{-p}\text{-set}\ d\ m$
and $snd\ ' snd\ ' set\ ps \subseteq set\ hs \times (set\ gs \cup set\ bs \cup set\ hs)$ **and** $unique\text{-idx}\ (gs\ @\ bs\ @\ hs)\ data$
and $is\text{-Groebner}\text{-basis}\ (fst\ ' set\ gs)$
and $\bigwedge p' q'. (p', q') \in snd\ ' set\ (apply\text{-ncrit}\ crit\ data\ gs\ bs\ hs\ ps) \implies$
 $fst\ p' \neq 0 \implies fst\ q' \neq 0 \implies crit\text{-pair}\text{-cbelow}\text{-on}\ d\ m\ (fst\ ' B)\ (fst\ p')\ (fst\ q')$
and $\bigwedge p' q'. p' \in set\ gs \cup set\ bs \implies q' \in set\ gs \cup set\ bs \implies fst\ p' \neq 0 \implies fst\ q' \neq 0 \implies$
 $crit\text{-pair}\text{-cbelow}\text{-on}\ d\ m\ (fst\ ' B)\ (fst\ p')\ (fst\ q')$
assumes $(ic, q\text{-in-}bs, p, q) \in set\ ps$ **and** $fst\ p \neq 0$ **and** $fst\ q \neq 0$
and $q\text{-in-}bs \implies (q \in set\ gs \cup set\ bs)$
shows $crit\text{-pair}\text{-cbelow}\text{-on}\ d\ m\ (fst\ ' B)\ (fst\ p)\ (fst\ q)$
 $\langle proof \rangle$

6.3.5 Applying Criteria to Old Pairs

definition $apply\text{-ocrit} :: ('t, 'b, 'c, 'd)\ ocritT \Rightarrow (nat \times 'd) \Rightarrow ('t, 'b, 'c)\ pdata\ list$
 \Rightarrow
 $(bool \times ('t, 'b, 'c)\ pdata\text{-pair})\ list \Rightarrow ('t, 'b, 'c)\ pdata\text{-pair}$
 $list \Rightarrow$
 $('t, 'b, 'c)\ pdata\text{-pair}\ list$

where $apply\text{-ocrit}\ crit\ data\ hs\ ps' ps = (let\ c = crit\ data\ hs\ ps'\ in\ [(p, q) \leftarrow ps.\ \neg\ c\ p\ q])$

lemma $set\text{-apply}\text{-ocrit}$:

$set\ (apply\text{-ocrit}\ crit\ data\ hs\ ps'\ ps) = \{(p, q) \mid p\ q.\ (p, q) \in set\ ps \wedge \neg\ crit\ data\ hs\ ps'\ p\ q\}$
 $\langle proof \rangle$

corollary *set-apply-ocrit-iff*:

$(p, q) \in \text{set } (\text{apply-ocrit crit data hs ps}' ps) \iff ((p, q) \in \text{set } ps \wedge \neg \text{crit data hs ps}' p q)$
 ⟨proof⟩

lemma *apply-ocrit-connectible*:

assumes *ocrit-spec crit and dickson-grading d and set hs* $\subseteq B$ **and** *fst ' B* \subseteq *dgrad-p-set d m*
and *unique-idx* $(p \# q \# \text{hs} @ (\text{map } (\text{fst} \circ \text{snd}) \text{ps}') @ (\text{map } (\text{snd} \circ \text{snd}) \text{ps}'))$
data
and $\bigwedge p' q'. (p', q') \in \text{snd } ' \text{set } \text{ps}' \implies \text{fst } p' \neq 0 \implies \text{fst } q' \neq 0 \implies$
 $\text{crit-pair-cbelow-on } d m (\text{fst } ' B) (\text{fst } p') (\text{fst } q')$
assumes $p \in B$ **and** $q \in B$ **and** $\text{fst } p \neq 0$ **and** $\text{fst } q \neq 0$
and $(p, q) \in \text{set } ps$ **and** $(p, q) \notin \text{set } (\text{apply-ocrit crit data hs ps}' ps)$
shows *crit-pair-cbelow-on d m (fst ' B) (fst p) (fst q)*
 ⟨proof⟩

6.3.6 Creating Final List of Pairs

context

fixes $np::('t, 'b::\text{field}, 'c, 'd) \text{np}T$
and $icrit::('t, 'b, 'c, 'd) \text{icrit}T$
and $ncrit::('t, 'b, 'c, 'd) \text{ncrit}T$
and $ocrit::('t, 'b, 'c, 'd) \text{ocrit}T$
and $\text{comb}::('t, 'b, 'c) \text{pdata-pair list} \Rightarrow ('t, 'b, 'c) \text{pdata-pair list} \Rightarrow ('t, 'b, 'c)$
pdata-pair list
begin

definition *add-pairs* $:: ('t, 'b, 'c, 'd) \text{ap}T$

where *add-pairs gs bs ps hs data* =
 $(\text{let } ps1 = \text{apply-ncrit ncrit data gs bs hs } (\text{apply-icrit icrit data gs bs hs } (np \text{ gs bs hs data}));$
 $ps2 = \text{apply-ocrit ocrit data hs ps1 ps in comb } (\text{map snd } [x \leftarrow ps1 . \neg \text{fst } x]) \text{ ps2})$

lemma *set-add-pairs*:

assumes $\bigwedge xs \text{ ys. set } (\text{comb } xs \text{ ys}) = \text{set } xs \cup \text{set } ys$
assumes $ps1 = \text{apply-ncrit ncrit data gs bs hs } (\text{apply-icrit icrit data gs bs hs } (np \text{ gs bs hs data}))$
shows $\text{set } (\text{add-pairs } gs \text{ bs } ps \text{ hs data}) =$
 $\{(p, q) \mid p \text{ q. } (\text{False}, p, q) \in \text{set } ps1 \vee ((p, q) \in \text{set } ps \wedge \neg \text{ocrit data hs ps1 } p \text{ q})\}$
 ⟨proof⟩

lemma *set-add-pairs-iff*:

assumes $\bigwedge xs \text{ ys. set } (\text{comb } xs \text{ ys}) = \text{set } xs \cup \text{set } ys$
assumes $ps1 = \text{apply-ncrit ncrit data gs bs hs } (\text{apply-icrit icrit data gs bs hs } (np \text{ gs bs hs data}))$

shows $((p, q) \in \text{set } (\text{add-pairs } gs \ bs \ ps \ hs \ data)) \longleftrightarrow$
 $((\text{False}, p, q) \in \text{set } ps1 \vee ((p, q) \in \text{set } ps \wedge \neg \text{ocrit } data \ hs \ ps1 \ p \ q))$
 $\langle \text{proof} \rangle$

lemma *ap-spec-add-pairs*:

assumes *np-spec np* **and** *icrit-spec icrit* **and** *ncrit-spec ncrit* **and** *ocrit-spec ocrit*
and $\bigwedge xs \ ys. \text{set } (\text{comb } xs \ ys) = \text{set } xs \cup \text{set } ys$
shows *ap-spec add-pairs*
 $\langle \text{proof} \rangle$

end

abbreviation *add-pairs-canon* \equiv

add-pairs (new-pairs-sorted canon-pair-order) component-crit chain-ncrit chain-ocrit
canon-pair-comb

lemma *ap-spec-add-pairs-canon*: *ap-spec add-pairs-canon*

$\langle \text{proof} \rangle$

6.4 Suitable Instances of the *completion* Parameter

definition *rcp-spec* :: $('t, 'b::\text{field}, 'c, 'd) \text{ compl}T \Rightarrow \text{bool}$

where *rcp-spec rcp* \longleftrightarrow
 $(\forall gs \ bs \ ps \ sps \ data.$
 $0 \notin \text{fst } ' \text{set } (\text{fst } (\text{rcp } gs \ bs \ ps \ sps \ data)) \wedge$
 $(\forall h \ b. h \in \text{set } (\text{fst } (\text{rcp } gs \ bs \ ps \ sps \ data)) \longrightarrow b \in \text{set } gs \cup \text{set } bs \longrightarrow$
 $\text{fst } b \neq 0 \longrightarrow$
 $\neg \text{lt } (\text{fst } b) \ \text{adds}_t \ \text{lt } (\text{fst } h)) \wedge$
 $(\forall d. \text{dickson-grading } d \longrightarrow$
 $\text{dgrad-p-set-le } d \ (\text{fst } ' \text{set } (\text{fst } (\text{rcp } gs \ bs \ ps \ sps \ data))) \ (\text{args-to-set}$
 $(gs, bs, sps))) \wedge$
 $\text{component-of-term } ' \text{Keys } (\text{fst } ' (\text{set } (\text{fst } (\text{rcp } gs \ bs \ ps \ sps \ data)))) \subseteq$
 $\text{component-of-term } ' \text{Keys } (\text{args-to-set } (gs, bs, sps)) \wedge$
 $(\text{is-Groebner-basis } (\text{fst } ' \text{set } gs) \longrightarrow \text{unique-idx } (gs @ bs) \ data \longrightarrow$
 $(\text{fst } ' \text{set } (\text{fst } (\text{rcp } gs \ bs \ ps \ sps \ data))) \subseteq \text{pmdl } (\text{args-to-set } (gs, bs, sps)))$
 \wedge
 $(\forall (p, q) \in \text{set } sps. \text{set } sps \subseteq \text{set } bs \times (\text{set } gs \cup \text{set } bs) \longrightarrow$
 $(\text{red } (\text{fst } ' (\text{set } gs \cup \text{set } bs) \cup \text{fst } ' \text{set } (\text{fst } (\text{rcp } gs \ bs \ ps \ sps \ data))))^{**}$
 $(\text{spoly } (\text{fst } p) \ (\text{fst } q)) \ 0))))$

Informally, *rcp-spec rcp* expresses that, for suitable *gs*, *bs* and *sps*, the value of *rcp gs bs ps sps*

- is a list consisting exclusively of non-zero polynomials contained in the module generated by $\text{set } bs \cup \text{set } gs$, whose leading terms are not divisible by the leading term of any non-zero $b \in \text{set } bs$, and
- contains sufficiently many new polynomials such that all S-polynomials originating from *sps* can be reduced to 0 modulo the enlarged list of polynomials.

lemma *rcp-specI*:

assumes $\bigwedge gs\ bs\ ps\ sps\ data. 0 \notin fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data))$
assumes $\bigwedge gs\ bs\ ps\ sps\ h\ b\ data. h \in set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data)) \implies b \in set\ gs \cup set\ bs \implies fst\ b \neq 0 \implies$
 $\neg lt\ (fst\ b)\ adds_t\ lt\ (fst\ h)$
assumes $\bigwedge gs\ bs\ ps\ sps\ d\ data. dickson-grading\ d \implies$
 $dgrad-p-set-le\ d\ (fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data)))\ (args-to-set\ (gs,\ bs,\ sps))$
assumes $\bigwedge gs\ bs\ ps\ sps\ data. component-of-term\ 'Keys\ (fst\ '(set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data)))) \subseteq$
 $component-of-term\ 'Keys\ (args-to-set\ (gs,\ bs,\ sps))$
assumes $\bigwedge gs\ bs\ ps\ sps\ data. is-Groebner-basis\ (fst\ 'set\ gs) \implies unique-idx\ (gs\ @\ bs)\ data \implies$
 $(fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data))) \subseteq pmdl\ (args-to-set\ (gs,\ bs,\ sps))$
 \wedge
 $(\forall (p, q) \in set\ sps. set\ sps \subseteq set\ bs \times (set\ gs \cup set\ bs) \longrightarrow$
 $(red\ (fst\ '(set\ gs \cup set\ bs) \cup fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data))))^{**}$
 $(spoly\ (fst\ p)\ (fst\ q))\ 0))$
shows *rcp-spec rcp*
 $\langle proof \rangle$

lemma *rcp-specD1*:

assumes *rcp-spec rcp*
shows $0 \notin fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data))$
 $\langle proof \rangle$

lemma *rcp-specD2*:

assumes *rcp-spec rcp*
and $h \in set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data))$ **and** $b \in set\ gs \cup set\ bs$ **and** $fst\ b \neq 0$
shows $\neg lt\ (fst\ b)\ adds_t\ lt\ (fst\ h)$
 $\langle proof \rangle$

lemma *rcp-specD3*:

assumes *rcp-spec rcp* **and** *dickson-grading d*
shows $dgrad-p-set-le\ d\ (fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data)))\ (args-to-set\ (gs,\ bs,\ sps))$
 $\langle proof \rangle$

lemma *rcp-specD4*:

assumes *rcp-spec rcp*
shows $component-of-term\ 'Keys\ (fst\ '(set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data)))) \subseteq$
 $component-of-term\ 'Keys\ (args-to-set\ (gs,\ bs,\ sps))$
 $\langle proof \rangle$

lemma *rcp-specD5*:

assumes *rcp-spec rcp* **and** *is-Groebner-basis (fst 'set gs)* **and** *unique-idx (gs @ bs) data*
shows $fst\ 'set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data)) \subseteq pmdl\ (args-to-set\ (gs,\ bs,\ sps))$
 $\langle proof \rangle$

lemma *rcp-specD6*:
assumes *rcp-spec rcp* **and** *is-Groebner-basis (fst ' set gs)* **and** *unique-idx (gs @ bs) data*
and $set\ sps \subseteq set\ bs \times (set\ gs \cup set\ bs)$
and $(p, q) \in set\ sps$
shows $(red\ (fst\ ' (set\ gs \cup set\ bs) \cup fst\ ' set\ (fst\ (rcp\ gs\ bs\ ps\ sps\ data))))^{**}$
 $(spoly\ (fst\ p)\ (fst\ q))\ 0$
 $\langle proof \rangle$

lemma *compl-struct-rcp*:
assumes *rcp-spec rcp*
shows *compl-struct rcp*
 $\langle proof \rangle$

lemma *compl-pmdl-rcp*:
assumes *rcp-spec rcp*
shows *compl-pmdl rcp*
 $\langle proof \rangle$

lemma *compl-conn-rcp*:
assumes *rcp-spec rcp*
shows *compl-conn rcp*
 $\langle proof \rangle$

end

6.5 Suitable Instances of the *add-basis* Parameter

definition *add-basis-naive* :: $('a, 'b, 'c, 'd)\ abT$
where *add-basis-naive gs bs ns data = bs @ ns*

lemma *ab-spec-add-basis-naive*: *ab-spec add-basis-naive*
 $\langle proof \rangle$

definition *add-basis-sorted* :: $(nat \times 'd \Rightarrow ('a, 'b, 'c)\ pdata \Rightarrow ('a, 'b, 'c)\ pdata \Rightarrow bool) \Rightarrow ('a, 'b, 'c, 'd)\ abT$
where *add-basis-sorted rel gs bs ns data = merge-wrt (rel data) bs ns*

lemma *ab-spec-add-basis-sorted*: *ab-spec (add-basis-sorted rel)*
 $\langle proof \rangle$

definition *card-keys* :: $('a \Rightarrow_0 'b::zero) \Rightarrow nat$
where *card-keys = card o keys*

definition (**in** *ordered-term*) *canon-basis-order* :: $'d \Rightarrow ('t, 'b::zero, 'c)\ pdata \Rightarrow ('t, 'b, 'c)\ pdata \Rightarrow bool$
where *canon-basis-order data p q \longleftrightarrow*
(let cp = card-keys (fst p); cq = card-keys (fst q) in

$$cp < cq \vee (cp = cq \wedge lt (fst p) \prec_t lt (fst q))$$

abbreviation (in *ordered-term*) *add-basis-canon* \equiv *add-basis-sorted canon-basis-order*

6.6 Special Case: Scalar Polynomials

context *gd-powerprod*

begin

lemma *remdups-map-component-of-term-punit*:

remdups (map (λ -. ()) (punit.Keys-to-list (map fst bs))) =
(if ($\forall b \in \text{set } bs. \text{fst } b = 0$) then [] else [()])
 ⟨*proof*⟩

lemma *count-const-lt-components-punit* [code]:

punit.count-const-lt-components hs =
(if ($\exists h \in \text{set } hs. \text{punit.const-lt-component (fst } h) = \text{Some } ()$) then 1 else 0)
 ⟨*proof*⟩

lemma *count-rem-components-punit* [code]:

punit.count-rem-components bs =
(if ($\forall b \in \text{set } bs. \text{fst } b = 0$) then 0
else
if ($\exists b \in \text{set } bs. \text{fst } b \neq 0 \wedge \text{punit.const-lt-component (fst } b) = \text{Some } ()$) then
0 else 1)
 ⟨*proof*⟩

lemma *full-gb-punit* [code]:

punit.full-gb bs = (if ($\forall b \in \text{set } bs. \text{fst } b = 0$) then [] else [(1, 0, default)])
 ⟨*proof*⟩

abbreviation *add-pairs-punit-canon* \equiv

punit.add-pairs (punit.new-pairs-sorted punit.canon-pair-order) punit.product-crit
punit.chain-ncrit
punit.chain-ocrit punit.canon-pair-comb

lemma *ap-spec-add-pairs-punit-canon*: *punit.ap-spec add-pairs-punit-canon*

⟨*proof*⟩

end

end

7 Buchberger's Algorithm

theory *Buchberger*

imports *Algorithm-Schema*

begin

context *gd-term*
begin

7.1 Reduction

definition *trdsp*::('t \Rightarrow_0 'b) list \Rightarrow ('t, 'b, 'c) pdata-pair \Rightarrow ('t \Rightarrow_0 'b::field)
where *trdsp bs p* \equiv *trd bs (spoly (fst (fst p)) (fst (snd p)))*

lemma *trdsp-alt*: *trdsp bs (p, q) = trd bs (spoly (fst p) (fst q))*
<proof>

lemma *trdsp-in-pmdl*: *trdsp bs (p, q) \in pmdl (insert (fst p) (insert (fst q) (set bs)))*
<proof>

lemma *dgrad-p-set-le-trdsp*:
assumes *dickson-grading d*
shows *dgrad-p-set-le d {trdsp bs (p, q)} (insert (fst p) (insert (fst q) (set bs)))*
<proof>

lemma *components-trdsp-subset*:
component-of-term ' keys (trdsp bs (p, q)) \subseteq component-of-term ' Keys (insert (fst p) (insert (fst q) (set bs)))
<proof>

definition *gb-red-aux* :: ('t, 'b::field, 'c) pdata list \Rightarrow ('t, 'b, 'c) pdata-pair list \Rightarrow
('t \Rightarrow_0 'b) list
where *gb-red-aux bs ps =*
(*let bs' = map fst bs in*
filter ($\lambda h. h \neq 0$) (map (trdsp bs') ps)
)

Actually, *gb-red-aux* is only called on singleton lists.

lemma *set-gb-red-aux*: *set (gb-red-aux bs ps) = (trdsp (map fst bs)) ' set ps - {0}*
<proof>

lemma *in-set-gb-red-auxI*:
assumes $(p, q) \in \text{set } ps$ **and** $h = \text{trdsp (map fst bs) (p, q)}$ **and** $h \neq 0$
shows $h \in \text{set (gb-red-aux bs ps)}$
<proof>

lemma *in-set-gb-red-auxE*:
assumes $h \in \text{set (gb-red-aux bs ps)}$
obtains $p q$ **where** $(p, q) \in \text{set } ps$ **and** $h = \text{trdsp (map fst bs) (p, q)}$
<proof>

lemma *gb-red-aux-not-zero*: $0 \notin \text{set (gb-red-aux bs ps)}$
<proof>

lemma *gb-red-aux-irreducible*:

assumes $h \in \text{set } (gb\text{-red-}aux\ bs\ ps)$ **and** $b \in \text{set } bs$ **and** $\text{fst } b \neq 0$

shows $\neg \text{lt } (\text{fst } b) \text{ adds}_t \text{ lt } h$

<proof>

lemma *gb-red-aux-dgrad-p-set-le*:

assumes *dickson-grading* d

shows *dgrad-p-set-le* d ($\text{set } (gb\text{-red-}aux\ bs\ ps)$) (*args-to-set* ($[], bs, ps$))

<proof>

lemma *components-gb-red-aux-subset*:

component-of-term ' *Keys* ($\text{set } (gb\text{-red-}aux\ bs\ ps)$) \subseteq *component-of-term* ' *Keys*
(*args-to-set* ($[], bs, ps$))

<proof>

lemma *pmdl-gb-red-aux*: $\text{set } (gb\text{-red-}aux\ bs\ ps) \subseteq \text{pmdl } (\text{args-to-set } ([], bs, ps))$

<proof>

lemma *gb-red-aux-spoly-reducible*:

assumes $(p, q) \in \text{set } ps$

shows $(\text{red } (\text{fst } ' \text{set } bs \cup \text{set } (gb\text{-red-}aux\ bs\ ps)))^{**} (\text{spoly } (\text{fst } p) (\text{fst } q)) = 0$

<proof>

definition *gb-red* :: ('t, 'b::field, 'c::default, 'd) *complT*

where *gb-red* $gs\ bs\ ps\ sps\ data = (\text{map } (\lambda h. (h, \text{default}))) (gb\text{-red-}aux\ (gs\ @\ bs)\ sps), \text{snd } data)$

lemma *fst-set-fst-gb-red*: $\text{fst } ' \text{set } (\text{fst } (gb\text{-red } gs\ bs\ ps\ sps\ data)) = \text{set } (gb\text{-red-}aux\ (gs\ @\ bs)\ sps)$

<proof>

lemma *rcp-spec-gb-red*: *rcp-spec* *gb-red*

<proof>

lemmas *compl-struct-gb-red* = *compl-struct-rcp*[*OF* *rcp-spec-gb-red*]

lemmas *compl-pmdl-gb-red* = *compl-pmdl-rcp*[*OF* *rcp-spec-gb-red*]

lemmas *compl-conn-gb-red* = *compl-conn-rcp*[*OF* *rcp-spec-gb-red*]

7.2 Pair Selection

primrec *gb-sel* :: ('t, 'b::zero, 'c, 'd) *selT* **where**

gb-sel $gs\ bs\ []\ data = []$

gb-sel $gs\ bs\ (p \# ps)\ data = [p]$

lemma *sel-spec-gb-sel*: *sel-spec* *gb-sel*

<proof>

7.3 Buchberger's Algorithm

lemma *struct-spec-gb*: *struct-spec* *gb-sel* *add-pairs-canon* *add-basis-canon* *gb-red*

<proof>

definition *gb-aux* :: ('t, 'b, 'c) pdata list \Rightarrow nat \times nat \times 'd \Rightarrow ('t, 'b, 'c) pdata list \Rightarrow

('t, 'b, 'c) pdata-pair list \Rightarrow ('t, 'b::field, 'c::default) pdata list

where *gb-aux* = *gb-schema-aux gb-sel add-pairs-canon add-basis-canon gb-red*

lemmas *gb-aux-simps* [code] = *gb-schema-aux-simps*[OF *struct-spec-gb, folded gb-aux-def*]

definition *gb* :: ('t, 'b, 'c) pdata' list \Rightarrow 'd \Rightarrow ('t, 'b::field, 'c::default) pdata' list

where *gb* = *gb-schema-direct gb-sel add-pairs-canon add-basis-canon gb-red*

lemmas *gb-simps* [code] = *gb-schema-direct-def*[of *gb-sel add-pairs-canon add-basis-canon gb-red, folded gb-def gb-aux-def*]

lemmas *gb-isGB* = *gb-schema-direct-isGB*[OF *struct-spec-gb compl-conn-gb-red, folded gb-def*]

lemmas *gb-pmdl* = *gb-schema-direct-pmdl*[OF *struct-spec-gb compl-pmdl-gb-red, folded gb-def*]

7.3.1 Special Case: *punit*

lemma (in *gd-term*) *struct-spec-gb-punit*: *punit.struct-spec punit.gb-sel add-pairs-punit-canon punit.add-basis-canon punit.gb-red*

<proof>

definition *gb-aux-punit* :: ('a, 'b, 'c) pdata list \Rightarrow nat \times nat \times 'd \Rightarrow ('a, 'b, 'c) pdata list \Rightarrow

('a, 'b, 'c) pdata-pair list \Rightarrow ('a, 'b::field, 'c::default) pdata list

where *gb-aux-punit* = *punit.gb-schema-aux punit.gb-sel add-pairs-punit-canon punit.add-basis-canon punit.gb-red*

lemmas *gb-aux-punit-simps* [code] = *punit.gb-schema-aux-simps*[OF *struct-spec-gb-punit, folded gb-aux-punit-def*]

definition *gb-punit* :: ('a, 'b, 'c) pdata' list \Rightarrow 'd \Rightarrow ('a, 'b::field, 'c::default) pdata' list

where *gb-punit* = *punit.gb-schema-direct punit.gb-sel add-pairs-punit-canon punit.add-basis-canon punit.gb-red*

lemmas *gb-punit-simps* [code] = *punit.gb-schema-direct-def*[of *punit.gb-sel add-pairs-punit-canon punit.add-basis-canon punit.gb-red, folded gb-punit-def gb-aux-punit-def*]

lemmas *gb-punit-isGB* = *punit.gb-schema-direct-isGB*[OF *struct-spec-gb-punit punit.compl-conn-gb-red, folded gb-punit-def*]

lemmas *gb-punit-pmdl* = *punit.gb-schema-direct-pmdl*[OF *struct-spec-gb-punit punit.compl-pmdl-gb-red,*

folded gb-punit-def]

end

end

8 Benchmark Problems for Computing Gröbner Bases

theory *Benchmarks*

imports *Polynomials.MPoly-Type-Class-OAlist*

begin

This theory defines various well-known benchmark problems for computing Gröbner bases. The actual tests of the different algorithms on these problems are contained in the theories whose names end with *-Examples*.

8.1 Cyclic

definition *cycl-pp* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow (\text{nat}, \text{nat}) \text{ pp}$

where *cycl-pp* $n\ d\ i = \text{sparse}_0 (\text{map } (\lambda k. (\text{modulo } (k + i)\ n, 1)) [0..<d])$

definition *cyclic* :: $(\text{nat}, \text{nat}) \text{ pp } \text{nat-term-order} \Rightarrow \text{nat} \Rightarrow ((\text{nat}, \text{nat}) \text{ pp} \Rightarrow_0 'a::\{\text{zero,one,uminus}\}) \text{ list}$

where *cyclic* $\text{to } n =$

($\text{let } xs = [0..<n]$ in
 $(\text{map } (\lambda d. \text{distr}_0 \text{ to } (\text{map } (\lambda i. (\text{cycl-pp } n\ d\ i, 1))\ xs)) [1..<n]) @$
 $[\text{distr}_0 \text{ to } [(\text{cycl-pp } n\ n\ 0, 1), (0, -1)]]$
 $)$

cyclic n is a system of n polynomials in n indeterminates, with maximum degree n .

8.2 Katsura

definition *katsura-poly* :: $(\text{nat}, \text{nat}) \text{ pp } \text{nat-term-order} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow ((\text{nat}, \text{nat}) \text{ pp} \Rightarrow_0 'a::\text{comm-ring-1})$

where *katsura-poly* $\text{to } n\ i =$

$\text{change-ord to } ((\sum j::\text{int}=-\text{int } n..<n + 1. \text{if } \text{abs } (i - j) \leq n \text{ then } V_0 (\text{nat } (\text{abs } j)) * V_0 (\text{nat } (\text{abs } (i - j))) \text{ else } 0) - V_0\ i)$

definition *katsura* :: $(\text{nat}, \text{nat}) \text{ pp } \text{nat-term-order} \Rightarrow \text{nat} \Rightarrow ((\text{nat}, \text{nat}) \text{ pp} \Rightarrow_0 'a::\text{comm-ring-1}) \text{ list}$

where *katsura* $\text{to } n =$

($\text{let } xs = [0..<n]$ in
 $(\text{distr}_0 \text{ to } ((\text{sparse}_0 [(0, 1)], 1) \# (\text{map } (\lambda i. (\text{sparse}_0 [(\text{Suc } i, 1)], 2))\ xs)) @ [(0, -1)]) \#$
 $(\text{map } (\text{katsura-poly to } n)\ xs)$

)

For $(1::'a) \leq n$, *katsura* n is a system of $n + 1$ polynomials in $n + 1$ indeterminates, with maximum degree 2.

8.3 Eco

definition *eco-poly* :: (nat, nat) pp nat-term-order \Rightarrow nat \Rightarrow nat \Rightarrow ((nat, nat) pp \Rightarrow_0 'a::comm-ring-1)

where *eco-poly to m i* =
 $\text{distr}_0 \text{ to } ((\text{sparse}_0 [(i, 1), (m, 1)], 1) \# \text{map } (\lambda j. (\text{sparse}_0 [(j, 1), (j + i + 1, 1), (m, 1)], 1)) [0..<m - i - 1])$

definition *eco* :: (nat, nat) pp nat-term-order \Rightarrow nat \Rightarrow ((nat, nat) pp \Rightarrow_0 'a::comm-ring-1) list

where *eco to n* =
 $(\text{let } m = n - 1 \text{ in}$
 $(\text{distr}_0 \text{ to } ((\text{map } (\lambda j. (\text{sparse}_0 [(j, 1)], 1)) [0..<m]) @ [(0, 1)])) \#$
 $(\text{distr}_0 \text{ to } [(\text{sparse}_0 [(m-1, 1), (m, 1)], 1), (0, - \text{of-nat } m)]) \#$
 $(\text{rev } (\text{map } (\text{eco-poly to } m) [0..<m-1]))$
 $)$

For $(2::'a) \leq n$, *eco* n is a system of n polynomials in n indeterminates, with maximum degree 3.

8.4 Noon

definition *noon-poly* :: (nat, nat) pp nat-term-order \Rightarrow nat \Rightarrow nat \Rightarrow ((nat, nat) pp \Rightarrow_0 'a::comm-ring-1)

where *noon-poly to n i* =
 $(\text{let } \text{ten} = \text{of-nat } 10; \text{eleven} = - \text{of-nat } 11 \text{ in}$
 $\text{distr}_0 \text{ to } ((\text{map } (\lambda j. \text{if } j = i \text{ then } (\text{sparse}_0 [(i, 1)], \text{eleven}) \text{ else } (\text{sparse}_0 [(j, 2), (i, 1)], \text{ten}))) [0..<n]) @$
 $[(0, \text{ten})])$

definition *noon* :: (nat, nat) pp nat-term-order \Rightarrow nat \Rightarrow ((nat, nat) pp \Rightarrow_0 'a::comm-ring-1) list

where *noon to n* = $(\text{noon-poly to } n \ 1) \# (\text{noon-poly to } n \ 0) \# (\text{map } (\text{noon-poly to } n) [2..<n])$

For $(2::'a) \leq n$, *noon* n is a system of n polynomials in n indeterminates, with maximum degree 3.

end

9 Code Equations Related to the Computation of Gröbner Bases

theory *Algorithm-Schema-Impl*

```

imports Algorithm-Schema Benchmarks
begin

lemma card-keys-MP-oalist [code]: card-keys (MP-oalist xs) = length (fst (list-of-oalist-ntm xs))
  <proof>

end

theory Code-Target-Rat
  imports Complex-Main HOL-Library.Code-Target-Numeral
begin

Mapping type rat to type "Rat.rat" in Isabelle/ML. Serialization for other
target languages will be provided in the future.

context includes integer.lifting begin

lift-definition rat-of-integer :: integer  $\Rightarrow$  rat is Rat.of-int <proof>

lift-definition quotient-of' :: rat  $\Rightarrow$  integer  $\times$  integer is quotient-of <proof>

lemma [code]: Rat.of-int (int-of-integer x) = rat-of-integer x
  <proof>

lemma [code-unfold]: quotient-of = ( $\lambda x.$  map-prod int-of-integer int-of-integer (quotient-of' x))
  <proof>

end

code-printing
type-constructor rat  $\rightarrow$ 
  (SML) Rat.rat |
constant plus :: rat  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Rat.add |
constant minus :: rat  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Rat.add ((-)) (Rat.neg ((-))) |
constant times :: rat  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Rat.mult |
constant inverse :: rat  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Rat.inv |
constant divide :: rat  $\Rightarrow$  -  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Rat.mult ((-)) (Rat.inv ((-))) |
constant rat-of-integer :: integer  $\Rightarrow$  rat  $\rightarrow$ 
  (SML) Rat.of'-int |
constant abs :: rat  $\Rightarrow$  -  $\rightarrow$ 
  (SML) Rat.abs |
constant 0 :: rat  $\rightarrow$ 

```

```

    (SML) !(Rat.make (0, 1)) |
constant 1 :: rat →
    (SML) !(Rat.make (1, 1)) |
constant uminus :: rat ⇒ rat →
    (SML) Rat.neg |
constant HOL.equal :: rat ⇒ - →
    (SML) !((- : Rat.rat) = -) |
constant quotient-of' →
    (SML) Rat.dest

```

end

10 Sample Computations with Buchberger's Algorithm

```

theory Buchberger-Examples
  imports Buchberger Algorithm-Schema-Impl Code-Target-Rat
begin

```

lemma (in *gd-term*) *compute-trd-aux* [code]:

```

  trd-aux fs p r =
    (if is-zero p then
      r
    else
      case find-adds fs (lt p) of
        None ⇒ trd-aux fs (tail p) (plus-monomial-less r (lc p) (lt p))
      | Some f ⇒ trd-aux fs (tail p - monom-mult (lc p / lc f) (lp p - lp f) (tail
f)) r
    )
  ⟨proof⟩

```

10.1 Scalar Polynomials

global-interpretation *punit'*: *gd-powerprod ord-pp-punit cmp-term ord-pp-strict-punit cmp-term*

```

rewrites punit.adds-term = (adds)
and punit.pp-of-term = (λx. x)
and punit.component-of-term = (λ-. ())
and punit.monom-mult = monom-mult-punit
and punit.mult-scalar = mult-scalar-punit
and punit'.punit.min-term = min-term-punit
and punit'.punit.lt = lt-punit cmp-term
and punit'.punit.lc = lc-punit cmp-term
and punit'.punit.tail = tail-punit cmp-term
and punit'.punit.ord-p = ord-p-punit cmp-term

```

and *punit'.punit.ord-strict-p* = *ord-strict-p-punit cmp-term*
for *cmp-term* :: ('a::nat, 'b::{nat,add-wellorder}) pp nat-term-order

defines *find-adds-punit* = *punit'.punit.find-adds*
and *trd-aux-punit* = *punit'.punit.trd-aux*
and *trd-punit* = *punit'.punit.trd*
and *spoly-punit* = *punit'.punit.spoly*
and *count-const-lt-components-punit* = *punit'.punit.count-const-lt-components*
and *count-rem-components-punit* = *punit'.punit.count-rem-components*
and *const-lt-component-punit* = *punit'.punit.const-lt-component*
and *full-gb-punit* = *punit'.punit.full-gb*
and *add-pairs-single-sorted-punit* = *punit'.punit.add-pairs-single-sorted*
and *add-pairs-punit* = *punit'.punit.add-pairs*
and *canon-pair-order-aux-punit* = *punit'.punit.canon-pair-order-aux*
and *canon-basis-order-punit* = *punit'.punit.canon-basis-order*
and *new-pairs-sorted-punit* = *punit'.punit.new-pairs-sorted*
and *product-crit-punit* = *punit'.punit.product-crit*
and *chain-ncrit-punit* = *punit'.punit.chain-ncrit*
and *chain-ocrit-punit* = *punit'.punit.chain-ocrit*
and *apply-icrit-punit* = *punit'.punit.apply-icrit*
and *apply-ncrit-punit* = *punit'.punit.apply-ncrit*
and *apply-ocrit-punit* = *punit'.punit.apply-ocrit*
and *trdsp-punit* = *punit'.punit.trdsp*
and *gb-sel-punit* = *punit'.punit.gb-sel*
and *gb-red-aux-punit* = *punit'.punit.gb-red-aux*
and *gb-red-punit* = *punit'.punit.gb-red*
and *gb-aux-punit* = *punit'.punit.gb-aux-punit*
and *gb-punit* = *punit'.punit.gb-punit* — Faster, because incorporates product criterion.
 ⟨*proof*⟩

lemma *compute-spoly-punit* [code]:
spoly-punit to p q = (let *t1* = *lt-punit to p*; *t2* = *lt-punit to q*; *l* = *lcs t1 t2* in
 (*monom-mult-punit* (*1 / lc-punit to p*) (*l - t1*) *p*) - (*monom-mult-punit*
 (*1 / lc-punit to q*) (*l - t2*) *q*)
 ⟨*proof*⟩

lemma *compute-trd-punit* [code]: *trd-punit to fs p* = *trd-aux-punit to fs p* (*change-ord to 0*)
 ⟨*proof*⟩

experiment begin interpretation *trivariate₀-rat* ⟨*proof*⟩

lemma
lt-punit DRLEX ($X^2 * Z^3 + 3 * X^2 * Y$) = *sparse₀* [(0, 2), (2, 3)]
 ⟨*proof*⟩

lemma
lc-punit DRLEX ($X^2 * Z^3 + 3 * X^2 * Y$) = 1

<proof>

lemma

tail-punit DRLEX $(X^2 * Z^3 + 3 * X^2 * Y) = 3 * X^2 * Y$

<proof>

lemma

ord-strict-p-punit DRLEX $(X^2 * Z^4 - 2 * Y^3 * Z^2) (X^2 * Z^7 + 2 * Y^3 * Z^2)$

<proof>

lemma

trd-punit DRLEX $[Y^2 * Z + 2 * Y * Z^3] (X^2 * Z^4 - 2 * Y^3 * Z^3) =$

$X^2 * Z^4 + Y^4 * Z$

<proof>

lemma

spoly-punit DRLEX $(X^2 * Z^4 - 2 * Y^3 * Z^2) (Y^2 * Z + 2 * Z^3) =$

$-2 * Y^3 * Z^2 - (C_0 (1 / 2)) * X^2 * Y^2 * Z^2$

<proof>

lemma

gb-punit DRLEX

$[$
 $(X^2 * Z^4 - 2 * Y^3 * Z^2, ()),$
 $(Y^2 * Z + 2 * Z^3, ())$

$] () =$

$[$
 $(-2 * Y^3 * Z^2 - (C_0 (1 / 2)) * X^2 * Y^2 * Z^2, ()),$
 $(X^2 * Z^4 - 2 * Y^3 * Z^2, ()),$

$(Y^2 * Z + 2 * Z^3, ()),$

$(-(C_0 (1 / 2)) * X^2 * Y^4 * Z - 2 * Y^5 * Z, ())$

$]$

<proof>

lemma

gb-punit DRLEX

$[$
 $(X^2 * Z^2 - Y, ()),$
 $(Y^2 * Z - 1, ())$

$] () =$

$[$
 $(-(Y^3) + X^2 * Z, ()),$
 $(X^2 * Z^2 - Y, ()),$

$(Y^2 * Z - 1, ())$

$]$

<proof>

```

lemma
  gb-punit DRLEX
  [
    ( $X^3 - X * Y * Z^2$ , ()),
    ( $Y^2 * Z - 1$ , ())
  ] () =
  [
    ( $-(X^3 * Y) + X * Z$ , ()),
    ( $X^3 - X * Y * Z^2$ , ()),
    ( $Y^2 * Z - 1$ , ()),
    ( $-(X * Z^3) + X^5$ , ())
  ]
  ⟨proof⟩

```

```

lemma
  gb-punit DRLEX
  [
    ( $X^2 + Y^2 + Z^2 - 1$ , ()),
    ( $X * Y - Z - 1$ , ()),
    ( $Y^2 + X$ , ()),
    ( $Z^2 + X$ , ())
  ] () =
  [
    (1, ())
  ]
  ⟨proof⟩

```

end

```

value [code] length (gb-punit DRLEX (map (λp. (p, ())) ((katsura DRLEX 2)::(-
⇒0 rat) list)) ())

```

```

value [code] length (gb-punit DRLEX (map (λp. (p, ())) ((cyclic DRLEX 5)::(-
⇒0 rat) list)) ())

```

10.2 Vector Polynomials

We must define the following four constants outside the global interpretation, since otherwise their types are too general.

```

definition splus-pprod :: ('a::nat, 'b::nat) pp ⇒ -
  where splus-pprod = pprod.splus

```

```

definition monom-mult-pprod :: 'c::semiring-0 ⇒ ('a::nat, 'b::nat) pp ⇒ -
  where monom-mult-pprod = pprod.monom-mult

```

```

definition mult-scalar-pprod :: (('a::nat, 'b::nat) pp ⇒0 'c::semiring-0) ⇒ -
  where mult-scalar-pprod = pprod.mult-scalar

```

```

definition adds-term-pprod :: (('a::nat, 'b::nat) pp × -) ⇒ -

```

where $\text{adds-term-pprod} = \text{pprod.adds-term}$

global-interpretation pprod' : $\text{gd-nat-term } \lambda x::('a, 'b) \text{ pp} \times 'c. x \lambda x. x \text{ cmp-term}$
 rewrites $\text{pprod.pp-of-term} = \text{fst}$
 and $\text{pprod.component-of-term} = \text{snd}$
 and $\text{pprod.splus} = \text{splus-pprod}$
 and $\text{pprod.monom-mult} = \text{monom-mult-pprod}$
 and $\text{pprod.mult-scalar} = \text{mult-scalar-pprod}$
 and $\text{pprod.adds-term} = \text{adds-term-pprod}$
 for $\text{cmp-term} :: (('a::\text{nat}, 'b::\text{nat}) \text{ pp} \times 'c::\{\text{nat}, \text{the-min}\}) \text{ nat-term-order}$
 defines $\text{shift-map-keys-pprod} = \text{pprod}'.\text{shift-map-keys}$
 and $\text{min-term-pprod} = \text{pprod}'.\text{min-term}$
 and $\text{lt-pprod} = \text{pprod}'.\text{lt}$
 and $\text{lc-pprod} = \text{pprod}'.\text{lc}$
 and $\text{tail-pprod} = \text{pprod}'.\text{tail}$
 and $\text{comp-opt-p-pprod} = \text{pprod}'.\text{comp-opt-p}$
 and $\text{ord-p-pprod} = \text{pprod}'.\text{ord-p}$
 and $\text{ord-strict-p-pprod} = \text{pprod}'.\text{ord-strict-p}$
 and $\text{find-adds-pprod} = \text{pprod}'.\text{find-adds}$
 and $\text{trd-aux-pprod} = \text{pprod}'.\text{trd-aux}$
 and $\text{trd-pprod} = \text{pprod}'.\text{trd}$
 and $\text{spoly-pprod} = \text{pprod}'.\text{spoly}$
 and $\text{count-const-lt-components-pprod} = \text{pprod}'.\text{count-const-lt-components}$
 and $\text{count-rem-components-pprod} = \text{pprod}'.\text{count-rem-components}$
 and $\text{const-lt-component-pprod} = \text{pprod}'.\text{const-lt-component}$
 and $\text{full-gb-pprod} = \text{pprod}'.\text{full-gb}$
 and $\text{keys-to-list-pprod} = \text{pprod}'.\text{keys-to-list}$
 and $\text{Keys-to-list-pprod} = \text{pprod}'.\text{Keys-to-list}$
 and $\text{add-pairs-single-sorted-pprod} = \text{pprod}'.\text{add-pairs-single-sorted}$
 and $\text{add-pairs-pprod} = \text{pprod}'.\text{add-pairs}$
 and $\text{canon-pair-order-aux-pprod} = \text{pprod}'.\text{canon-pair-order-aux}$
 and $\text{canon-basis-order-pprod} = \text{pprod}'.\text{canon-basis-order}$
 and $\text{new-pairs-sorted-pprod} = \text{pprod}'.\text{new-pairs-sorted}$
 and $\text{component-crit-pprod} = \text{pprod}'.\text{component-crit}$
 and $\text{chain-ncrit-pprod} = \text{pprod}'.\text{chain-ncrit}$
 and $\text{chain-ocrit-pprod} = \text{pprod}'.\text{chain-ocrit}$
 and $\text{apply-icrit-pprod} = \text{pprod}'.\text{apply-icrit}$
 and $\text{apply-ncrit-pprod} = \text{pprod}'.\text{apply-ncrit}$
 and $\text{apply-ocrit-pprod} = \text{pprod}'.\text{apply-ocrit}$
 and $\text{trdsp-pprod} = \text{pprod}'.\text{trdsp}$
 and $\text{gb-sel-pprod} = \text{pprod}'.\text{gb-sel}$
 and $\text{gb-red-aux-pprod} = \text{pprod}'.\text{gb-red-aux}$
 and $\text{gb-red-pprod} = \text{pprod}'.\text{gb-red}$
 and $\text{gb-aux-pprod} = \text{pprod}'.\text{gb-aux}$
 and $\text{gb-pprod} = \text{pprod}'.\text{gb}$
 $\langle \text{proof} \rangle$

lemma $\text{compute-adds-term-pprod}$ [code]:
 $\text{adds-term-pprod } u \ v = (\text{snd } u = \text{snd } v \wedge \text{adds-pp-add-linorder } (\text{fst } u) (\text{fst } v))$

<proof>

lemma *compute-splus-pprod* [code]: *splus-pprod* $t (s, i) = (t + s, i)$
<proof>

lemma *compute-shift-map-keys-pprod* [code abstract]:
list-of-oalist-ntm (*shift-map-keys-pprod* $t f xs$) = *map-raw* $(\lambda(k, v). (splus-pprod$
 $t k, f v))$ (*list-of-oalist-ntm* xs)
<proof>

lemma *compute-trd-pprod* [code]: *trd-pprod* $to fs p = trd-aux-pprod$ $to fs p$ (*change-ord*
to 0)
<proof>

lemmas [code] = *conversep-iff*

definition $Vec_0 :: nat \Rightarrow (('a, nat) pp \Rightarrow_0 'b) \Rightarrow (('a::nat, nat) pp \times nat) \Rightarrow_0$
 $'b::semiring-1$ **where**
 $Vec_0 i p = mult-scalar-pprod p (Poly-Mapping.single (0, i) 1)$

experiment begin interpretation *trivariate₀-rat* *<proof>*

lemma
ord-p-pprod (*POT DRLEX*) ($Vec_0 1 (X^2 * Z) + Vec_0 0 (2 * Y \wedge 3 * Z^2)$) (Vec_0
 $1 (X^2 * Z^2 + 2 * Y \wedge 3 * Z^2)$)
<proof>

lemma
tail-pprod (*POT DRLEX*) ($Vec_0 1 (X^2 * Z) + Vec_0 0 (2 * Y \wedge 3 * Z^2)$) = Vec_0
 $0 (2 * Y \wedge 3 * Z^2)$
<proof>

lemma
lt-pprod (*POT DRLEX*) ($Vec_0 1 (X^2 * Z) + Vec_0 0 (2 * Y \wedge 3 * Z^2)$) = (*sparse₀*
 $[(0, 2), (2, 1)], 1$)
<proof>

lemma
keys ($Vec_0 0 (X^2 * Z \wedge 3) + Vec_0 1 (2 * Y \wedge 3 * Z^2)$) =
 $\{(sparse_0 [(0, 2), (2, 3)], 0), (sparse_0 [(1, 3), (2, 2)], 1)\}$
<proof>

lemma
keys ($Vec_0 0 (X^2 * Z \wedge 3) + Vec_0 2 (2 * Y \wedge 3 * Z^2)$) =
 $\{(sparse_0 [(0, 2), (2, 3)], 0), (sparse_0 [(1, 3), (2, 2)], 2)\}$
<proof>

lemma
 $Vec_0 1 (X^2 * Z \wedge 7 + 2 * Y \wedge 3 * Z^2) + Vec_0 3 (X^2 * Z \wedge 4) + Vec_0 1 (- 2$

* $Y^3 * Z^2 =$
 $Vec_0 1 (X^2 * Z^7) + Vec_0 3 (X^2 * Z^4)$
 ⟨proof⟩

lemma
 $lookup (Vec_0 0 (X^2 * Z^7) + Vec_0 1 (2 * Y^3 * Z^2 + 2)) (sparse_0 [(0, 2), (2, 7)], 0) = 1$
 ⟨proof⟩

lemma
 $lookup (Vec_0 0 (X^2 * Z^7) + Vec_0 1 (2 * Y^3 * Z^2 + 2)) (sparse_0 [(0, 2), (2, 7)], 1) = 0$
 ⟨proof⟩

lemma
 $Vec_0 0 (0 * X^2 * Z^7) + Vec_0 1 (0 * Y^3 * Z^2) = 0$
 ⟨proof⟩

lemma
 $monom-mult-pprod 3 (sparse_0 [(1, 2::nat)]) (Vec_0 0 (X^2 * Z) + Vec_0 1 (2 * Y^3 * Z^2)) =$
 $Vec_0 0 (3 * Y^2 * Z * X^2) + Vec_0 1 (6 * Y^5 * Z^2)$
 ⟨proof⟩

lemma
 $trd-pprod DRLEX [Vec_0 0 (Y^2 * Z + 2 * Y * Z^3)] (Vec_0 0 (X^2 * Z^4 - 2 * Y^3 * Z^3)) =$
 $Vec_0 0 (X^2 * Z^4 + Y^4 * Z)$
 ⟨proof⟩

lemma
 $length (gb-pprod (POT DRLEX$
 $[$
 $(Vec_0 0 (X^2 * Z^4 - 2 * Y^3 * Z^2), ()),$
 $(Vec_0 0 (Y^2 * Z + 2 * Z^3), ())$
 $]) = 4$
 ⟨proof⟩

end

end

11 Further Properties of Multivariate Polynomials

theory *More-MPoly-Type-Class*
imports *Polynomials.MPoly-Type-Class-Ordered General*
begin

Some further general properties of (ordered) multivariate polynomials needed for Gröbner bases. This theory is an extension of *Polynomials.MPoly-Type-Class-Ordered*.

11.1 Modules and Linear Hulls

context *module*
begin

lemma *span-listE*:

assumes $p \in \text{span } (\text{set } bs)$

obtains qs **where** $\text{length } qs = \text{length } bs$ **and** $p = \text{sum-list } (\text{map2 } (*s) \text{ } qs \text{ } bs)$

<proof>

lemma *span-listI*: $\text{sum-list } (\text{map2 } (*s) \text{ } qs \text{ } bs) \in \text{span } (\text{set } bs)$

<proof>

end

lemma (**in** *term-powerprod*) *monomial-1-in-pmdlI*:

assumes $(f::- \Rightarrow_0 'b::\text{field}) \in \text{pmdl } F$ **and** $\text{keys } f = \{t\}$

shows *monomial 1 t* $\in \text{pmdl } F$

<proof>

11.2 Ordered Polynomials

context *ordered-term*
begin

11.2.1 Sets of Leading Terms and -Coefficients

definition *lt-set* :: $(t, 'b::\text{zero}) \text{ poly-mapping set} \Rightarrow 't \text{ set}$ **where**

$\text{lt-set } F = \text{lt } ' (F - \{0\})$

definition *lc-set* :: $(t, 'b::\text{zero}) \text{ poly-mapping set} \Rightarrow 'b \text{ set}$ **where**

$\text{lc-set } F = \text{lc } ' (F - \{0\})$

lemma *lt-setI*:

assumes $f \in F$ **and** $f \neq 0$

shows $\text{lt } f \in \text{lt-set } F$

<proof>

lemma *lt-setE*:

assumes $t \in \text{lt-set } F$

obtains f **where** $f \in F$ **and** $f \neq 0$ **and** $\text{lt } f = t$

<proof>

lemma *lt-set-iff*:

shows $t \in \text{lt-set } F \iff (\exists f \in F. f \neq 0 \wedge \text{lt } f = t)$

<proof>

lemma *lc-setI*:

assumes $f \in F$ and $f \neq 0$

shows $lc\ f \in lc\text{-set } F$

$\langle proof \rangle$

lemma *lc-setE*:

assumes $c \in lc\text{-set } F$

obtains f where $f \in F$ and $f \neq 0$ and $lc\ f = c$

$\langle proof \rangle$

lemma *lc-set-iff*:

shows $c \in lc\text{-set } F \iff (\exists f \in F. f \neq 0 \wedge lc\ f = c)$

$\langle proof \rangle$

lemma *lc-set-nonzero*:

shows $0 \notin lc\text{-set } F$

$\langle proof \rangle$

lemma *lt-sum-distinct-eq-Max*:

assumes *finite* I and $sum\ p\ I \neq 0$

and $\bigwedge i1\ i2. i1 \in I \implies i2 \in I \implies p\ i1 \neq 0 \implies p\ i2 \neq 0 \implies lt\ (p\ i1) = lt\ (p\ i2) \implies i1 = i2$

shows $lt\ (sum\ p\ I) = ord\text{-term-lin.Max}\ (lt\text{-set}\ (p\ 'I))$

$\langle proof \rangle$

lemma *lt-sum-distinct-in-lt-set*:

assumes *finite* I and $sum\ p\ I \neq 0$

and $\bigwedge i1\ i2. i1 \in I \implies i2 \in I \implies p\ i1 \neq 0 \implies p\ i2 \neq 0 \implies lt\ (p\ i1) = lt\ (p\ i2) \implies i1 = i2$

shows $lt\ (sum\ p\ I) \in lt\text{-set}\ (p\ 'I)$

$\langle proof \rangle$

11.2.2 Monicity

definition *monic* :: $(t \Rightarrow_0 'b) \Rightarrow (t \Rightarrow_0 'b::field)$ where

$monic\ p = monom\text{-mult}\ (1 / lc\ p)\ 0\ p$

definition *is-monic-set* :: $(t \Rightarrow_0 'b::field)$ set $\Rightarrow bool$ where

$is\text{-monic-set}\ B \equiv (\forall b \in B. b \neq 0 \longrightarrow lc\ b = 1)$

lemma *lookup-monic*: $lookup\ (monic\ p)\ v = (lookup\ p\ v) / lc\ p$

$\langle proof \rangle$

lemma *lookup-monic-lt*:

assumes $p \neq 0$

shows $lookup\ (monic\ p)\ (lt\ p) = 1$

$\langle proof \rangle$

lemma *monic-0* [*simp*]: $\text{monic } 0 = 0$
<proof>

lemma *monic-0-iff*: $(\text{monic } p = 0) \longleftrightarrow (p = 0)$
<proof>

lemma *keys-monic* [*simp*]: $\text{keys } (\text{monic } p) = \text{keys } p$
<proof>

lemma *lt-monic* [*simp*]: $\text{lt } (\text{monic } p) = \text{lt } p$
<proof>

lemma *lc-monic*:
 assumes $p \neq 0$
 shows $\text{lc } (\text{monic } p) = 1$
<proof>

lemma *mult-lc-monic*:
 assumes $p \neq 0$
 shows $\text{monom-mult } (\text{lc } p) \ 0 \ (\text{monic } p) = p$ (**is** $?q = p$)
<proof>

lemma *is-monic-setI*:
 assumes $\bigwedge b. b \in B \implies b \neq 0 \implies \text{lc } b = 1$
 shows *is-monic-set* B
<proof>

lemma *is-monic-setD*:
 assumes *is-monic-set* B **and** $b \in B$ **and** $b \neq 0$
 shows $\text{lc } b = 1$
<proof>

lemma *Keys-image-monic* [*simp*]: $\text{Keys } (\text{monic } ` A) = \text{Keys } A$
<proof>

lemma *image-monic-is-monic-set*: *is-monic-set* $(\text{monic } ` A)$
<proof>

lemma *pmdl-image-monic* [*simp*]: $\text{pmdl } (\text{monic } ` B) = \text{pmdl } B$
<proof>

end

end

12 Auto-reducing Lists of Polynomials

theory *Auto-Reduction*

imports *Reduction More-MPoly-Type-Class*

begin

12.1 Reduction and Monic Sets

context *ordered-term*

begin

lemma *is-red-monic*: $is-red\ B\ (monic\ p) \longleftrightarrow is-red\ B\ p$
<proof>

lemma *red-image-monic* [*simp*]: $red\ (monic\ 'B) = red\ B$
<proof>

lemma *is-red-image-monic* [*simp*]: $is-red\ (monic\ 'B)\ p \longleftrightarrow is-red\ B\ p$
<proof>

12.2 Minimal Bases and Auto-reduced Bases

definition *is-auto-reduced* :: $(t \Rightarrow_0 'b::field)\ set \Rightarrow bool$ **where**
 $is-auto-reduced\ B \equiv (\forall b \in B. \neg is-red\ (B - \{b\})\ b)$

definition *is-minimal-basis* :: $(t \Rightarrow_0 'b::zero)\ set \Rightarrow bool$ **where**
 $is-minimal-basis\ B \longleftrightarrow (0 \notin B \wedge (\forall p\ q. p \in B \longrightarrow q \in B \longrightarrow p \neq q \longrightarrow \neg lt\ p\ adds_t\ lt\ q))$

lemma *is-auto-reducedD*:

assumes *is-auto-reduced* B **and** $b \in B$

shows $\neg is-red\ (B - \{b\})\ b$

<proof>

The converse of the following lemma is only true if B is minimal!

lemma *image-monic-is-auto-reduced*:

assumes *is-auto-reduced* B

shows *is-auto-reduced* $(monic\ 'B)$

<proof>

lemma *is-minimal-basisI*:

assumes $\bigwedge p. p \in B \Longrightarrow p \neq 0$ **and** $\bigwedge p\ q. p \in B \Longrightarrow q \in B \Longrightarrow p \neq q \Longrightarrow \neg lt\ p\ adds_t\ lt\ q$

shows *is-minimal-basis* B

<proof>

lemma *is-minimal-basisD1*:

assumes *is-minimal-basis* B **and** $p \in B$

shows $p \neq 0$

<proof>

lemma *is-minimal-basisD2*:

assumes *is-minimal-basis* B **and** $p \in B$ **and** $q \in B$ **and** $p \neq q$

shows $\neg lt\ p\ adds_t\ lt\ q$
<proof>

lemma *is-minimal-basisD3*:
assumes *is-minimal-basis* B **and** $p \in B$ **and** $q \in B$ **and** $p \neq q$
shows $\neg lt\ q\ adds_t\ lt\ p$
<proof>

lemma *is-minimal-basis-subset*:
assumes *is-minimal-basis* B **and** $A \subseteq B$
shows *is-minimal-basis* A
<proof>

lemma *nadds-red*:
assumes *nadds*: $\bigwedge q. q \in B \implies \neg lt\ q\ adds_t\ lt\ p$ **and** *red*: *red* $B\ p\ r$
shows $r \neq 0 \wedge lt\ r = lt\ p$
<proof>

lemma *nadds-red-nonzero*:
assumes *nadds*: $\bigwedge q. q \in B \implies \neg lt\ q\ adds_t\ lt\ p$ **and** *red* $B\ p\ r$
shows $r \neq 0$
<proof>

lemma *nadds-red-lt*:
assumes *nadds*: $\bigwedge q. q \in B \implies \neg lt\ q\ adds_t\ lt\ p$ **and** *red* $B\ p\ r$
shows $lt\ r = lt\ p$
<proof>

lemma *nadds-red-rtrancl-lt*:
assumes *nadds*: $\bigwedge q. q \in B \implies \neg lt\ q\ adds_t\ lt\ p$ **and** *rtrancl*: $(red\ B)^{**}\ p\ r$
shows $lt\ r = lt\ p$
<proof>

lemma *nadds-red-rtrancl-nonzero*:
assumes *nadds*: $\bigwedge q. q \in B \implies \neg lt\ q\ adds_t\ lt\ p$ **and** $p \neq 0$ **and** *rtrancl*: $(red\ B)^{**}\ p\ r$
shows $r \neq 0$
<proof>

lemma *minimal-basis-red-rtrancl-nonzero*:
assumes *is-minimal-basis* B **and** $p \in B$ **and** $(red\ (B - \{p\}))^{**}\ p\ r$
shows $r \neq 0$
<proof>

lemma *minimal-basis-red-rtrancl-lt*:
assumes *is-minimal-basis* B **and** $p \in B$ **and** $(red\ (B - \{p\}))^{**}\ p\ r$
shows $lt\ r = lt\ p$
<proof>

lemma *is-minimal-basis-replace*:

assumes *major*: *is-minimal-basis* B **and** $p \in B$ **and** *red*: $(\text{red } (B - \{p\}))^{**} p r$

shows *is-minimal-basis* $(\text{insert } r (B - \{p\}))$

<proof>

12.3 Computing Minimal Bases

definition *comp-min-basis* :: $(t \Rightarrow_0 'b)$ list $\Rightarrow (t \Rightarrow_0 'b::\text{zero})$ list **where**

comp-min-basis $xs = \text{filter-min } (\lambda x y. \text{lt } x \text{ adds}_t \text{ lt } y) (\text{filter } (\lambda x. x \neq 0) xs)$

lemma *comp-min-basis-subset'*: $\text{set } (\text{comp-min-basis } xs) \subseteq \{x \in \text{set } xs. x \neq 0\}$

<proof>

lemma *comp-min-basis-subset*: $\text{set } (\text{comp-min-basis } xs) \subseteq \text{set } xs$

<proof>

lemma *comp-min-basis-nonzero*: $p \in \text{set } (\text{comp-min-basis } xs) \Longrightarrow p \neq 0$

<proof>

lemma *comp-min-basis-adds*:

assumes $p \in \text{set } xs$ **and** $p \neq 0$

obtains q **where** $q \in \text{set } (\text{comp-min-basis } xs)$ **and** $\text{lt } q \text{ adds}_t \text{ lt } p$

<proof>

lemma *comp-min-basis-is-red*:

assumes *is-red* $(\text{set } xs)$ f

shows *is-red* $(\text{set } (\text{comp-min-basis } xs))$ f

<proof>

lemma *comp-min-basis-nadds*:

assumes $p \in \text{set } (\text{comp-min-basis } xs)$ **and** $q \in \text{set } (\text{comp-min-basis } xs)$ **and** $p \neq$

q

shows $\neg \text{lt } q \text{ adds}_t \text{ lt } p$

<proof>

lemma *comp-min-basis-is-minimal-basis*: *is-minimal-basis* $(\text{set } (\text{comp-min-basis } xs))$

<proof>

lemma *comp-min-basis-distinct*: *distinct* $(\text{comp-min-basis } xs)$

<proof>

end

12.4 Auto-Reduction

context *gd-term*

begin

lemma *is-minimal-basis-trd-is-minimal-basis*:

assumes *is-minimal-basis* $(\text{set } (x \# xs))$ **and** $x \notin \text{set } xs$

shows *is-minimal-basis* (set ((trd xs x) # xs))
 ⟨proof⟩

lemma *is-minimal-basis-trd-distinct*:

assumes *min*: *is-minimal-basis* (set (x # xs)) **and** *dist*: *distinct* (x # xs)
shows *distinct* ((trd xs x) # xs)
 ⟨proof⟩

primrec *comp-red-basis-aux* :: ('t \Rightarrow_0 'b) list \Rightarrow ('t \Rightarrow_0 'b) list \Rightarrow ('t \Rightarrow_0 'b::field) list **where**

comp-red-basis-aux-base: *comp-red-basis-aux* Nil ys = ys |
comp-red-basis-aux-rec: *comp-red-basis-aux* (x # xs) ys = *comp-red-basis-aux* xs
 ((trd (xs @ ys) x) # ys)

lemma *subset-comp-red-basis-aux*: set ys \subseteq set (*comp-red-basis-aux* xs ys)
 ⟨proof⟩

lemma *comp-red-basis-aux-nonzero*:

assumes *is-minimal-basis* (set (xs @ ys)) **and** *distinct* (xs @ ys) **and** $p \in$ set
 (*comp-red-basis-aux* xs ys)
shows $p \neq 0$
 ⟨proof⟩

lemma *comp-red-basis-aux-lt*:

assumes *is-minimal-basis* (set (xs @ ys)) **and** *distinct* (xs @ ys)
shows $lt \text{ ' set (xs @ ys) = } lt \text{ ' set (comp-red-basis-aux xs ys)}$
 ⟨proof⟩

lemma *comp-red-basis-aux-pmdl*:

assumes *is-minimal-basis* (set (xs @ ys)) **and** *distinct* (xs @ ys)
shows $pmdl$ (set (*comp-red-basis-aux* xs ys)) \subseteq $pmdl$ (set (xs @ ys))
 ⟨proof⟩

lemma *comp-red-basis-aux-irred*:

assumes *is-minimal-basis* (set (xs @ ys)) **and** *distinct* (xs @ ys)
and $\bigwedge y. y \in$ set ys $\implies \neg$ *is-red* (set (xs @ ys) - {y}) y
and $p \in$ set (*comp-red-basis-aux* xs ys)
shows \neg *is-red* (set (*comp-red-basis-aux* xs ys) - {p}) p
 ⟨proof⟩

lemma *comp-red-basis-aux-dgrad-p-set-le*:

assumes *dickson-grading* d
shows *dgrad-p-set-le* d (set (*comp-red-basis-aux* xs ys)) (set xs \cup set ys)
 ⟨proof⟩

definition *comp-red-basis* :: ('t \Rightarrow_0 'b) list \Rightarrow ('t \Rightarrow_0 'b::field) list

where *comp-red-basis* xs = *comp-red-basis-aux* (*comp-min-basis* xs) \square

lemma *comp-red-basis-nonzero*:

assumes $p \in \text{set } (\text{comp-red-basis } xs)$
shows $p \neq 0$
 $\langle \text{proof} \rangle$

lemma *pmdl-comp-red-basis-subset*: $\text{pmdl } (\text{set } (\text{comp-red-basis } xs)) \subseteq \text{pmdl } (\text{set } xs)$
 $\langle \text{proof} \rangle$

lemma *comp-red-basis-adds*:
assumes $p \in \text{set } xs$ **and** $p \neq 0$
obtains q **where** $q \in \text{set } (\text{comp-red-basis } xs)$ **and** $\text{lt } q \text{ adds}_t \text{ lt } p$
 $\langle \text{proof} \rangle$

lemma *comp-red-basis-lt*:
assumes $p \in \text{set } (\text{comp-red-basis } xs)$
obtains q **where** $q \in \text{set } xs$ **and** $q \neq 0$ **and** $\text{lt } q = \text{lt } p$
 $\langle \text{proof} \rangle$

lemma *comp-red-basis-is-red*: $\text{is-red } (\text{set } (\text{comp-red-basis } xs)) \iff \text{is-red } (\text{set } xs)$
 f
 $\langle \text{proof} \rangle$

lemma *comp-red-basis-is-auto-reduced*: $\text{is-auto-reduced } (\text{set } (\text{comp-red-basis } xs))$
 $\langle \text{proof} \rangle$

lemma *comp-red-basis-dgrad-p-set-le*:
assumes *dickson-grading* d
shows *dgrad-p-set-le* d $(\text{set } (\text{comp-red-basis } xs))$ $(\text{set } xs)$
 $\langle \text{proof} \rangle$

12.5 Auto-Reduction and Monicity

definition *comp-red-monic-basis* :: $(t \Rightarrow_0 b)$ list \Rightarrow $(t \Rightarrow_0 b::\text{field})$ list **where**
 $\text{comp-red-monic-basis } xs = \text{map } \text{monic } (\text{comp-red-basis } xs)$

lemma *set-comp-red-monic-basis*: $\text{set } (\text{comp-red-monic-basis } xs) = \text{monic } \text{' } (\text{set } (\text{comp-red-basis } xs))$
 $\langle \text{proof} \rangle$

lemma *comp-red-monic-basis-nonzero*:
assumes $p \in \text{set } (\text{comp-red-monic-basis } xs)$
shows $p \neq 0$
 $\langle \text{proof} \rangle$

lemma *comp-red-monic-basis-is-monic-set*: $\text{is-monic-set } (\text{set } (\text{comp-red-monic-basis } xs))$
 $\langle \text{proof} \rangle$

lemma *pmdl-comp-red-monic-basis-subset*: $\text{pmdl } (\text{set } (\text{comp-red-monic-basis } xs))$

\subseteq *pmdl* (set *xs*)
<proof>

lemma *comp-red-monic-basis-is-auto-reduced: is-auto-reduced* (set (comp-red-monic-basis *xs*))
<proof>

lemma *comp-red-monic-basis-dgrad-p-set-le:*
assumes *dickson-grading d*
shows *dgrad-p-set-le d* (set (comp-red-monic-basis *xs*)) (set *xs*)
<proof>

end

end

13 Reduced Gröbner Bases

theory *Reduced-GB*
imports *Groebner-Bases Auto-Reduction*
begin

lemma (in *gd-term*) *GB-image-monic: is-Groebner-basis* (monic ' *G*) \longleftrightarrow *is-Groebner-basis* *G*
<proof>

13.1 Definition and Uniqueness of Reduced Gröbner Bases

context *ordered-term*
begin

definition *is-reduced-GB* :: ('*t* \Rightarrow_0 '*b*::field) set \Rightarrow bool **where**
is-reduced-GB B \equiv *is-Groebner-basis B* \wedge *is-auto-reduced B* \wedge *is-monic-set B* \wedge
0 \notin *B*

lemma *reduced-GB-D1:*
assumes *is-reduced-GB G*
shows *is-Groebner-basis G*
<proof>

lemma *reduced-GB-D2:*
assumes *is-reduced-GB G*
shows *is-auto-reduced G*
<proof>

lemma *reduced-GB-D3:*
assumes *is-reduced-GB G*
shows *is-monic-set G*
<proof>

lemma *reduced-GB-D4*:
assumes *is-reduced-GB* G **and** $g \in G$
shows $g \neq 0$
 \langle *proof* \rangle

lemma *reduced-GB-lc*:
assumes *major: is-reduced-GB* G **and** $g \in G$
shows $lc\ g = 1$
 \langle *proof* \rangle

end

context *gd-term*
begin

lemma *is-reduced-GB-subsetI*:
assumes *Ared: is-reduced-GB* A **and** *BGB: is-Groebner-basis* B **and** *Bmon: is-monic-set* B
and $*$: $\bigwedge a\ b. a \in A \implies b \in B \implies a \neq 0 \implies b \neq 0 \implies a - b \neq 0 \implies lt\ (a - b) \in keys\ b \implies lt\ (a - b) \prec_t\ lt\ b \implies False$
and *id-eq: pmdl* $A = pmdl\ B$
shows $A \subseteq B$
 \langle *proof* \rangle

lemma *is-reduced-GB-unique'*:
assumes *Ared: is-reduced-GB* A **and** *Bred: is-reduced-GB* B **and** *id-eq: pmdl* $A = pmdl\ B$
shows $A \subseteq B$
 \langle *proof* \rangle

theorem *is-reduced-GB-unique*:
assumes *Ared: is-reduced-GB* A **and** *Bred: is-reduced-GB* B **and** *id-eq: pmdl* $A = pmdl\ B$
shows $A = B$
 \langle *proof* \rangle

13.2 Computing Reduced Gröbner Bases by Auto-Reduction

13.2.1 Minimal Bases

lemma *minimal-basis-is-reduced-GB*:
assumes *is-minimal-basis* B **and** *is-monic-set* B **and** *is-reduced-GB* G **and** $G \subseteq B$
and *pmdl* $B = pmdl\ G$
shows $B = G$
 \langle *proof* \rangle

13.2.2 Computing Minimal Bases

lemma *comp-min-basis-pmdl*:

assumes *is-Groebner-basis* (set *xs*)

shows $\text{pmdl} (\text{set} (\text{comp-min-basis } xs)) = \text{pmdl} (\text{set } xs)$ (**is** $\text{pmdl} (\text{set } ?ys) = -$)

<proof>

lemma *comp-min-basis-GB*:

assumes *is-Groebner-basis* (set *xs*)

shows *is-Groebner-basis* (set (comp-min-basis *xs*)) (**is** *is-Groebner-basis* (set *?ys*))

<proof>

13.2.3 Computing Reduced Bases

lemma *comp-red-basis-pmdl*:

assumes *is-Groebner-basis* (set *xs*)

shows $\text{pmdl} (\text{set} (\text{comp-red-basis } xs)) = \text{pmdl} (\text{set } xs)$

<proof>

lemma *comp-red-basis-GB*:

assumes *is-Groebner-basis* (set *xs*)

shows *is-Groebner-basis* (set (comp-red-basis *xs*))

<proof>

13.2.4 Computing Reduced Gröbner Bases

lemma *comp-red-monic-basis-pmdl*:

assumes *is-Groebner-basis* (set *xs*)

shows $\text{pmdl} (\text{set} (\text{comp-red-monic-basis } xs)) = \text{pmdl} (\text{set } xs)$

<proof>

lemma *comp-red-monic-basis-GB*:

assumes *is-Groebner-basis* (set *xs*)

shows *is-Groebner-basis* (set (comp-red-monic-basis *xs*))

<proof>

lemma *comp-red-monic-basis-is-reduced-GB*:

assumes *is-Groebner-basis* (set *xs*)

shows *is-reduced-GB* (set (comp-red-monic-basis *xs*))

<proof>

lemma *ex-finite-reduced-GB-dgrad-p-set*:

assumes *dickson-grading* *d* **and** *finite* (component-of-term ‘ *Keys F*’) **and** $F \subseteq \text{dgrad-p-set } d \ m$

obtains *G* **where** $G \subseteq \text{dgrad-p-set } d \ m$ **and** *finite* *G* **and** *is-reduced-GB* *G* **and** $\text{pmdl } G = \text{pmdl } F$

<proof>

theorem *ex-unique-reduced-GB-dgrad-p-set*:

assumes *dickson-grading* *d* **and** *finite* (component-of-term ‘ *Keys F*’) **and** $F \subseteq$

dgrad-p-set d m

shows $\exists! G. G \subseteq \text{dgrad-p-set } d \ m \wedge \text{finite } G \wedge \text{is-reduced-GB } G \wedge \text{pmdl } G = \text{pmdl } F$
<proof>

corollary *ex-unique-reduced-GB-dgrad-p-set'*:

assumes *dickson-grading d* **and** *finite (component-of-term ' Keys F)* **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows $\exists! G. \text{finite } G \wedge \text{is-reduced-GB } G \wedge \text{pmdl } G = \text{pmdl } F$
<proof>

definition *reduced-GB* :: $(t \Rightarrow_0 'b) \text{ set} \Rightarrow (t \Rightarrow_0 'b::\text{field}) \text{ set}$

where *reduced-GB B* = (THE $G. \text{finite } G \wedge \text{is-reduced-GB } G \wedge \text{pmdl } G = \text{pmdl } B$)

reduced-GB returns the unique reduced Gröbner basis of the given set, provided its Dickson grading is bounded. Combining *comp-red-monic-basis* with any function for computing Gröbner bases, e.g. *gb* from theory "Buchberger", makes *reduced-GB* computable.

lemma *finite-reduced-GB-dgrad-p-set*:

assumes *dickson-grading d* **and** *finite (component-of-term ' Keys F)* **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows *finite (reduced-GB F)*
<proof>

lemma *reduced-GB-is-reduced-GB-dgrad-p-set*:

assumes *dickson-grading d* **and** *finite (component-of-term ' Keys F)* **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows *is-reduced-GB (reduced-GB F)*
<proof>

lemma *reduced-GB-is-GB-dgrad-p-set*:

assumes *dickson-grading d* **and** *finite (component-of-term ' Keys F)* **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows *is-Groebner-basis (reduced-GB F)*
<proof>

lemma *reduced-GB-is-auto-reduced-dgrad-p-set*:

assumes *dickson-grading d* **and** *finite (component-of-term ' Keys F)* **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows *is-auto-reduced (reduced-GB F)*
<proof>

lemma *reduced-GB-is-monic-set-dgrad-p-set*:

assumes *dickson-grading d* **and** *finite (component-of-term ' Keys F)* **and** $F \subseteq \text{dgrad-p-set } d \ m$

shows *is-monic-set (reduced-GB F)*
<proof>

lemma *reduced-GB-nonzero-dgrad-p-set:*
assumes *dickson-grading d and finite (component-of-term ‘ Keys F) and $F \subseteq$*
dgrad-p-set d m
shows $0 \notin \text{reduced-GB } F$
<proof>

lemma *reduced-GB-pmdl-dgrad-p-set:*
assumes *dickson-grading d and finite (component-of-term ‘ Keys F) and $F \subseteq$*
dgrad-p-set d m
shows $\text{pmdl (reduced-GB } F) = \text{pmdl } F$
<proof>

lemma *reduced-GB-unique-dgrad-p-set:*
assumes *dickson-grading d and finite (component-of-term ‘ Keys F) and $F \subseteq$*
dgrad-p-set d m
and *is-reduced-GB G and $\text{pmdl } G = \text{pmdl } F$*
shows $\text{reduced-GB } F = G$
<proof>

lemma *reduced-GB-dgrad-p-set:*
assumes *dickson-grading d and finite (component-of-term ‘ Keys F) and $F \subseteq$*
dgrad-p-set d m
shows $\text{reduced-GB } F \subseteq \text{dgrad-p-set } d m$
<proof>

lemma *reduced-GB-unique:*
assumes *finite G and is-reduced-GB G and $\text{pmdl } G = \text{pmdl } F$*
shows $\text{reduced-GB } F = G$
<proof>

lemma *is-reduced-GB-empty: is-reduced-GB {}*
<proof>

lemma *is-reduced-GB-singleton: is-reduced-GB {f} \longleftrightarrow lc f = 1*
<proof>

lemma *reduced-GB-empty: reduced-GB {} = {}*
<proof>

lemma *reduced-GB-singleton: reduced-GB {f} = (if f = 0 then {} else {monic f})*
<proof>

lemma *ex-unique-reduced-GB-finite: finite F \implies ($\exists! G. \text{finite } G \wedge \text{is-reduced-GB}$*
 $G \wedge \text{pmdl } G = \text{pmdl } F$)
<proof>

lemma *finite-reduced-GB-finite: finite F \implies finite (reduced-GB F)*
<proof>

lemma *reduced-GB-is-reduced-GB-finite*: $\text{finite } F \implies \text{is-reduced-GB } (\text{reduced-GB } F)$

<proof>

lemma *reduced-GB-is-GB-finite*: $\text{finite } F \implies \text{is-Groebner-basis } (\text{reduced-GB } F)$

<proof>

lemma *reduced-GB-is-auto-reduced-finite*: $\text{finite } F \implies \text{is-auto-reduced } (\text{reduced-GB } F)$

<proof>

lemma *reduced-GB-is-monic-set-finite*: $\text{finite } F \implies \text{is-monic-set } (\text{reduced-GB } F)$

<proof>

lemma *reduced-GB-nonzero-finite*: $\text{finite } F \implies 0 \notin \text{reduced-GB } F$

<proof>

lemma *reduced-GB-pmdl-finite*: $\text{finite } F \implies \text{pmdl } (\text{reduced-GB } F) = \text{pmdl } F$

<proof>

lemma *reduced-GB-unique-finite*: $\text{finite } F \implies \text{is-reduced-GB } G \implies \text{pmdl } G = \text{pmdl } F \implies \text{reduced-GB } F = G$

<proof>

end

13.2.5 Properties of the Reduced Gröbner Basis of an Ideal

context *gd-powerprod*

begin

lemma *ideal-eq-UNIV-iff-reduced-GB-eq-one-dgrad-p-set*:

assumes *dickson-grading* d **and** $F \subseteq \text{punit.dgrad-p-set } d$

shows $\text{ideal } F = \text{UNIV} \iff \text{punit.reduced-GB } F = \{1\}$

<proof>

lemmas *ideal-eq-UNIV-iff-reduced-GB-eq-one-finite* =

ideal-eq-UNIV-iff-reduced-GB-eq-one-dgrad-p-set[*OF dickson-grading-dgrad-dummy*
punit.dgrad-p-set-exhaust-expl]

end

13.2.6 Context *od-term*

context *od-term*

begin

lemmas *ex-unique-reduced-GB* =

ex-unique-reduced-GB-dgrad-p-set'[*OF dickson-grading-zero - subset-dgrad-p-set-zero*]

lemmas *finite-reduced-GB* =

```

    finite-reduced-GB-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-is-reduced-GB =
    reduced-GB-is-reduced-GB-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-is-GB =
    reduced-GB-is-GB-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-is-auto-reduced =
    reduced-GB-is-auto-reduced-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-is-monic-set =
    reduced-GB-is-monic-set-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-nonzero =
    reduced-GB-nonzero-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-pmdl =
    reduced-GB-pmdl-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]
lemmas reduced-GB-unique =
    reduced-GB-unique-dgrad-p-set[OF dickson-grading-zero - subset-dgrad-p-set-zero]

end

end

```

14 Sample Computations of Reduced Gröbner Bases

```

theory Reduced-GB-Examples
  imports Buchberger Reduced-GB Polynomials.MPoly-Type-Class-OAlist Code-Target-Rat
begin

context gd-term
begin

definition rgb :: ('t  $\Rightarrow_0$  'b) list  $\Rightarrow$  ('t  $\Rightarrow_0$  'b::field) list
  where rgb bs = comp-red-monic-basis (map fst (gb (map ( $\lambda$ b. (b, ())) bs) ()))

definition rgb-punit :: ('a  $\Rightarrow_0$  'b) list  $\Rightarrow$  ('a  $\Rightarrow_0$  'b::field) list
  where rgb-punit bs = punit.comp-red-monic-basis (map fst (gb-punit (map ( $\lambda$ b.
  (b, ())) bs) ()))

lemma compute-trd-aux [code]:
  trd-aux fs p r =
    (if is-zero p then
     r
    else
     case find-adds fs (lt p) of
       None  $\Rightarrow$  trd-aux fs (tail p) (plus-monomial-less r (lc p) (lt p))
     | Some f  $\Rightarrow$  trd-aux fs (tail p - monom-mult (lc p / lc f) (lp p - lp f) (tail
f)) r
    )
  <proof>

end

```

We only consider scalar polynomials here, but vector-polynomials could be handled, too.

global-interpretation *punit'*: *gd-powerprod ord-pp-punit cmp-term ord-pp-strict-punit cmp-term*

```

rewrites punit.adds-term = (adds)
and punit.pp-of-term = ( $\lambda x. x$ )
and punit.component-of-term = ( $\lambda. ()$ )
and punit.monom-mult = monom-mult-punit
and punit.mult-scalar = mult-scalar-punit
and punit'.punit.min-term = min-term-punit
and punit'.punit.lt = lt-punit cmp-term
and punit'.punit.lc = lc-punit cmp-term
and punit'.punit.tail = tail-punit cmp-term
and punit'.punit.ord-p = ord-p-punit cmp-term
and punit'.punit.ord-strict-p = ord-strict-p-punit cmp-term
for cmp-term :: ('a::nat, 'b::{nat,add-wellorder}) pp nat-term-order

```

```

defines find-adds-punit = punit'.punit.find-adds
and trd-aux-punit = punit'.punit.trd-aux
and trd-punit = punit'.punit.trd
and spoly-punit = punit'.punit.spoly
and count-const-lt-components-punit = punit'.punit.count-const-lt-components
and count-rem-components-punit = punit'.punit.count-rem-components
and const-lt-component-punit = punit'.punit.const-lt-component
and full-gb-punit = punit'.punit.full-gb
and add-pairs-single-sorted-punit = punit'.punit.add-pairs-single-sorted
and add-pairs-punit = punit'.punit.add-pairs
and canon-pair-order-aux-punit = punit'.punit.canon-pair-order-aux
and canon-basis-order-punit = punit'.punit.canon-basis-order
and new-pairs-sorted-punit = punit'.punit.new-pairs-sorted
and product-crit-punit = punit'.punit.product-crit
and chain-ncrit-punit = punit'.punit.chain-ncrit
and chain-ocrit-punit = punit'.punit.chain-ocrit
and apply-icrit-punit = punit'.punit.apply-icrit
and apply-ncrit-punit = punit'.punit.apply-ncrit
and apply-ocrit-punit = punit'.punit.apply-ocrit
and trdsp-punit = punit'.punit.trdsp
and gb-sel-punit = punit'.punit.gb-sel
and gb-red-aux-punit = punit'.punit.gb-red-aux
and gb-red-punit = punit'.punit.gb-red
and gb-aux-punit = punit'.punit.gb-aux-punit
and gb-punit = punit'.punit.gb-punit — Faster, because incorporates product
criterion.
and comp-min-basis-punit = punit'.punit.comp-min-basis
and comp-red-basis-aux-punit = punit'.punit.comp-red-basis-aux
and comp-red-basis-punit = punit'.punit.comp-red-basis
and monic-punit = punit'.punit.monic
and comp-red-monic-basis-punit = punit'.punit.comp-red-monic-basis
and rgb-punit = punit'.punit.rgb-punit

```

<proof>

lemma *compute-spoly-punit* [code]:

spoly-punit to p q = (let t1 = lt-punit to p; t2 = lt-punit to q; l = lcs t1 t2 in
(monom-mult-punit (1 / lc-punit to p) (l - t1) p) - (monom-mult-punit
(1 / lc-punit to q) (l - t2) q))
<proof>

lemma *compute-trd-punit* [code]: *trd-punit to fs p = trd-aux-punit to fs p (change-ord*
to 0)

<proof>

experiment begin interpretation *trivariate₀-rat* *<proof>*

lemma

rgb-punit DRLEX

[
 $X^3 - X * Y * Z^2,$
 $Y^2 * Z - 1$
]
=
[
 $X^3 * Y - X * Z,$
 $-(X^3) + X * Y * Z^2,$
 $Y^2 * Z - 1,$
 $-(X * Z^3) + X^5$
]

<proof>

lemma

rgb-punit DRLEX

[
 $X^2 + Y^2 + Z^2 - 1,$
 $X * Y - Z - 1,$
 $Y^2 + X,$
 $Z^2 + X$
]
=
[
 1
]

<proof>

Note: The above computations have been cross-checked with Mathematica 11.1.

end

end

15 Macaulay Matrices

theory *Macaulay-Matrix*

imports *More-MPoly-Type-Class Jordan-Normal-Form.Gauss-Jordan-Elimination*
begin

We build upon vectors and matrices represented by dimension and characteristic function, because later on we need to quantify the dimensions of certain matrices existentially. This is not possible (at least not easily possible) with a type-based approach, as in HOL-Multivariate Analysis.

15.1 More about Vectors

lemma *vec-of-list-alt*: $vec\ of\ list\ xs = vec\ (length\ xs)\ (nth\ xs)$
<proof>

lemma *vec-cong*:

assumes $n = m$ **and** $\bigwedge i. i < m \implies f\ i = g\ i$
shows $vec\ n\ f = vec\ m\ g$
<proof>

lemma *scalar-prod-comm*:

assumes $dim\ vec\ v = dim\ vec\ w$
shows $v \cdot w = w \cdot (v :: 'a :: comm\ semiring\ 0\ vec)$
<proof>

lemma *vec-scalar-mult-fun*: $vec\ n\ (\lambda x. c * f\ x) = c \cdot_v\ vec\ n\ f$
<proof>

definition *mult-vec-mat* :: $'a\ vec \Rightarrow 'a :: semiring\ 0\ mat \Rightarrow 'a\ vec$ (**infixl** $_v*$ 70)
where $v \cdot_v\ A \equiv vec\ (dim\ col\ A)\ (\lambda j. v \cdot col\ A\ j)$

definition *resize-vec* :: $nat \Rightarrow 'a\ vec \Rightarrow 'a\ vec$
where $resize\ vec\ n\ v = vec\ n\ (vec\ index\ v)$

lemma *dim-resize-vec[simp]*: $dim\ vec\ (resize\ vec\ n\ v) = n$
<proof>

lemma *resize-vec-carrier*: $resize\ vec\ n\ v \in carrier\ vec\ n$
<proof>

lemma *resize-vec-dim[simp]*: $resize\ vec\ (dim\ vec\ v)\ v = v$
<proof>

lemma *resize-vec-index*:

assumes $i < n$
shows $resize\ vec\ n\ v\ \$\ i = v\ \$\ i$
<proof>

lemma *mult-mat-vec-resize*:

$v \cdot_v^* A = (\text{resize-vec } (\text{dim-row } A) \ v) \cdot_v^* A$
<proof>

lemma *assoc-mult-vec-mat*:

assumes $v \in \text{carrier-vec } n1$ **and** $A \in \text{carrier-mat } n1 \ n2$ **and** $B \in \text{carrier-mat } n2 \ n3$
shows $v \cdot_v^* (A * B) = (v \cdot_v^* A) \cdot_v^* B$
<proof>

lemma *mult-vec-mat-transpose*:

assumes $\text{dim-vec } v = \text{dim-row } A$
shows $v \cdot_v^* A = (\text{transpose-mat } A) \cdot_v (v :: 'a :: \text{comm-semiring-0 vec})$
<proof>

15.2 More about Matrices

definition *nzrows* :: $'a :: \text{zero mat} \Rightarrow 'a \text{ vec list}$

where $\text{nzrows } A = \text{filter } (\lambda r. r \neq 0_v (\text{dim-col } A)) (\text{rows } A)$

definition *row-space* :: $'a \text{ mat} \Rightarrow 'a :: \text{semiring-0 vec set}$

where $\text{row-space } A = (\lambda v. \text{mult-vec-mat } v \ A) \ ` (\text{carrier-vec } (\text{dim-row } A))$

definition *row-echelon* :: $'a \text{ mat} \Rightarrow 'a :: \text{field mat}$

where $\text{row-echelon } A = \text{fst } (\text{gauss-jordan } A \ (1_m (\text{dim-row } A)))$

15.2.1 nzrows

lemma *length-nzrows*: $\text{length } (\text{nzrows } A) \leq \text{dim-row } A$

<proof>

lemma *set-nzrows*: $\text{set } (\text{nzrows } A) = \text{set } (\text{rows } A) - \{0_v (\text{dim-col } A)\}$

<proof>

lemma *nzrows-nth-not-zero*:

assumes $i < \text{length } (\text{nzrows } A)$

shows $\text{nzrows } A \ ! \ i \neq 0_v (\text{dim-col } A)$

<proof>

15.2.2 row-space

lemma *row-spaceI*:

assumes $x = v \cdot_v^* A$

shows $x \in \text{row-space } A$

<proof>

lemma *row-spaceE*:

assumes $x \in \text{row-space } A$

obtains v **where** $v \in \text{carrier-vec } (\text{dim-row } A)$ **and** $x = v \cdot_v^* A$

<proof>

lemma *row-space-alt*: $\text{row-space } A = \text{range } (\lambda v. \text{mult-vec-mat } v \ A)$
(*proof*)

lemma *row-space-mult*:
 assumes $A \in \text{carrier-mat } nr \ nc$ **and** $B \in \text{carrier-mat } nr \ nr$
 shows $\text{row-space } (B * A) \subseteq \text{row-space } A$
(*proof*)

lemma *row-space-mult-unit*:
 assumes $P \in \text{Units } (\text{ring-mat } \text{TYPE}('a::\text{semiring-1}) \ (\text{dim-row } A) \ b)$
 shows $\text{row-space } (P * A) = \text{row-space } A$
(*proof*)

15.2.3 *row-echelon*

lemma *row-eq-zero-iff-pivot-fun*:
 assumes *pivot-fun* $A \ f \ (\text{dim-col } A)$ **and** $i < \text{dim-row } (A::'a::\text{zero-neq-one mat})$
 shows $(\text{row } A \ i = 0_v \ (\text{dim-col } A)) \longleftrightarrow (f \ i = \text{dim-col } A)$
(*proof*)

lemma *row-not-zero-iff-pivot-fun*:
 assumes *pivot-fun* $A \ f \ (\text{dim-col } A)$ **and** $i < \text{dim-row } (A::'a::\text{zero-neq-one mat})$
 shows $(\text{row } A \ i \neq 0_v \ (\text{dim-col } A)) \longleftrightarrow (f \ i < \text{dim-col } A)$
(*proof*)

lemma *pivot-fun-stabilizes*:
 assumes *pivot-fun* $A \ f \ nc$ **and** $i1 \leq i2$ **and** $i2 < \text{dim-row } A$ **and** $nc \leq f \ i1$
 shows $f \ i2 = nc$
(*proof*)

lemma *pivot-fun-mono-strict*:
 assumes *pivot-fun* $A \ f \ nc$ **and** $i1 < i2$ **and** $i2 < \text{dim-row } A$ **and** $f \ i1 < nc$
 shows $f \ i1 < f \ i2$
(*proof*)

lemma *pivot-fun-mono*:
 assumes *pivot-fun* $A \ f \ nc$ **and** $i1 \leq i2$ **and** $i2 < \text{dim-row } A$
 shows $f \ i1 \leq f \ i2$
(*proof*)

lemma *row-echelon-carrier*:
 assumes $A \in \text{carrier-mat } nr \ nc$
 shows $\text{row-echelon } A \in \text{carrier-mat } nr \ nc$
(*proof*)

lemma *dim-row-echelon[simp]*:
 shows $\text{dim-row } (\text{row-echelon } A) = \text{dim-row } A$ **and** $\text{dim-col } (\text{row-echelon } A) = \text{dim-col } A$

<proof>

lemma *row-echelon-transform*:

obtains P **where** $P \in \text{Units } (\text{ring-mat } \text{TYPE}('a::\text{field}) (\text{dim-row } A) b)$ **and**
 $\text{row-echelon } A = P * A$

<proof>

lemma *row-space-row-echelon[simp]*: $\text{row-space } (\text{row-echelon } A) = \text{row-space } A$

<proof>

lemma *row-echelon-pivot-fun*:

obtains f **where** $\text{pivot-fun } (\text{row-echelon } A) f (\text{dim-col } (\text{row-echelon } A))$

<proof>

lemma *distinct-nzrows-row-echelon*: $\text{distinct } (\text{nzrows } (\text{row-echelon } A))$

<proof>

15.3 Converting Between Polynomials and Macaulay Matrices

definition *poly-to-row* :: $'a \text{ list} \Rightarrow ('a \Rightarrow_0 'b::\text{zero}) \Rightarrow 'b \text{ vec}$ **where**
 $\text{poly-to-row } ts \ p = \text{vec-of-list } (\text{map } (\text{lookup } p) \ ts)$

definition *polys-to-mat* :: $'a \text{ list} \Rightarrow ('a \Rightarrow_0 'b::\text{zero}) \text{ list} \Rightarrow 'b \text{ mat}$ **where**
 $\text{polys-to-mat } ts \ ps = \text{mat-of-rows } (\text{length } ts) (\text{map } (\text{poly-to-row } ts) \ ps)$

definition *list-to-fun* :: $'a \text{ list} \Rightarrow ('b::\text{zero}) \text{ list} \Rightarrow 'a \Rightarrow 'b$ **where**
 $\text{list-to-fun } ts \ cs \ t = (\text{case map-of } (\text{zip } ts \ cs) \ t \text{ of } \text{Some } c \Rightarrow c \mid \text{None} \Rightarrow 0)$

definition *list-to-poly* :: $'a \text{ list} \Rightarrow 'b \text{ list} \Rightarrow ('a \Rightarrow_0 'b::\text{zero})$ **where**
 $\text{list-to-poly } ts \ cs = \text{Abs-poly-mapping } (\text{list-to-fun } ts \ cs)$

definition *row-to-poly* :: $'a \text{ list} \Rightarrow 'b \text{ vec} \Rightarrow ('a \Rightarrow_0 'b::\text{zero})$ **where**
 $\text{row-to-poly } ts \ r = \text{list-to-poly } ts (\text{list-of-vec } r)$

definition *mat-to-polys* :: $'a \text{ list} \Rightarrow 'b \text{ mat} \Rightarrow ('a \Rightarrow_0 'b::\text{zero}) \text{ list}$ **where**
 $\text{mat-to-polys } ts \ A = \text{map } (\text{row-to-poly } ts) (\text{rows } A)$

lemma *dim-poly-to-row*: $\text{dim-vec } (\text{poly-to-row } ts \ p) = \text{length } ts$
<proof>

lemma *poly-to-row-index*:

assumes $i < \text{length } ts$

shows $\text{poly-to-row } ts \ p \ \$ \ i = \text{lookup } p (ts \ ! \ i)$

<proof>

context *term-powerprod*

begin

lemma *poly-to-row-scalar-mult*:

assumes $keys\ p \subseteq set\ ts$

shows $row\text{-}to\text{-}poly\ ts\ (c \cdot_v\ (poly\text{-}to\text{-}row\ ts\ p)) = c \cdot p$

<proof>

lemma *poly-to-row-to-poly*:

assumes $keys\ p \subseteq set\ ts$

shows $row\text{-}to\text{-}poly\ ts\ (poly\text{-}to\text{-}row\ ts\ p) = (p::'t \Rightarrow_0\ 'b::semiring-1)$

<proof>

lemma *lookup-list-to-poly*: $lookup\ (list\text{-}to\text{-}poly\ ts\ cs) = list\text{-}to\text{-}fun\ ts\ cs$

<proof>

lemma *list-to-fun-Nil* [simp]: $list\text{-}to\text{-}fun\ []\ cs = 0$

<proof>

lemma *list-to-poly-Nil* [simp]: $list\text{-}to\text{-}poly\ []\ cs = 0$

<proof>

lemma *row-to-poly-Nil* [simp]: $row\text{-}to\text{-}poly\ []\ r = 0$

<proof>

lemma *lookup-row-to-poly*:

assumes *distinct* ts **and** $dim\text{-}vec\ r = length\ ts$ **and** $i < length\ ts$

shows $lookup\ (row\text{-}to\text{-}poly\ ts\ r)\ (ts\ !\ i) = r\ \$\ i$

<proof>

lemma *keys-row-to-poly*: $keys\ (row\text{-}to\text{-}poly\ ts\ r) \subseteq set\ ts$

<proof>

lemma *lookup-row-to-poly-not-zeroE*:

assumes $lookup\ (row\text{-}to\text{-}poly\ ts\ r)\ t \neq 0$

obtains i **where** $i < length\ ts$ **and** $t = ts\ !\ i$

<proof>

lemma *row-to-poly-zero* [simp]: $row\text{-}to\text{-}poly\ ts\ (0_v\ (length\ ts)) = (0::'t \Rightarrow_0\ 'b::zero)$

<proof>

lemma *row-to-poly-zeroD*:

assumes *distinct* ts **and** $dim\text{-}vec\ r = length\ ts$ **and** $row\text{-}to\text{-}poly\ ts\ r = 0$

shows $r = 0_v\ (length\ ts)$

<proof>

lemma *row-to-poly-inj*:

assumes *distinct* ts **and** $dim\text{-}vec\ r1 = length\ ts$ **and** $dim\text{-}vec\ r2 = length\ ts$

and $row\text{-}to\text{-}poly\ ts\ r1 = row\text{-}to\text{-}poly\ ts\ r2$

shows $r1 = r2$

<proof>

lemma *row-to-poly-vec-plus*:

assumes *distinct ts and length ts = n*

shows $\text{row-to-poly } ts \ (\text{vec } n \ (f1 + f2)) = \text{row-to-poly } ts \ (\text{vec } n \ f1) + \text{row-to-poly } ts \ (\text{vec } n \ f2)$

<proof>

lemma *row-to-poly-vec-sum*:

assumes *distinct ts and length ts = n*

shows $\text{row-to-poly } ts \ (\text{vec } n \ (\lambda j. \sum_{i \in I}. f \ i \ j)) = ((\sum_{i \in I}. \text{row-to-poly } ts \ (\text{vec } n \ (f \ i)))::'t \Rightarrow_0 \ 'b::\text{comm-monoid-add})$

<proof>

lemma *row-to-poly-smult*:

assumes *distinct ts and dim-vec r = length ts*

shows $\text{row-to-poly } ts \ (c \cdot_v \ r) = c \cdot (\text{row-to-poly } ts \ r)$

<proof>

lemma *poly-to-row-Nil [simp]*: $\text{poly-to-row } [] \ p = \text{vec } 0 \ f$

<proof>

lemma *polys-to-mat-Nil [simp]*: $\text{polys-to-mat } ts \ [] = \text{mat } 0 \ (\text{length } ts) \ f$

<proof>

lemma *dim-row-polys-to-mat [simp]*: $\text{dim-row } (\text{polys-to-mat } ts \ ps) = \text{length } ps$

<proof>

lemma *dim-col-polys-to-mat [simp]*: $\text{dim-col } (\text{polys-to-mat } ts \ ps) = \text{length } ts$

<proof>

lemma *polys-to-mat-index*:

assumes $i < \text{length } ps$ **and** $j < \text{length } ts$

shows $(\text{polys-to-mat } ts \ ps) \ \$\$ \ (i, j) = \text{lookup } (ps \ ! \ i) \ (ts \ ! \ j)$

<proof>

lemma *row-polys-to-mat*:

assumes $i < \text{length } ps$

shows $\text{row } (\text{polys-to-mat } ts \ ps) \ i = \text{poly-to-row } ts \ (ps \ ! \ i)$

<proof>

lemma *col-polys-to-mat*:

assumes $j < \text{length } ts$

shows $\text{col } (\text{polys-to-mat } ts \ ps) \ j = \text{vec-of-list } (\text{map } (\lambda p. \text{lookup } p \ (ts \ ! \ j)) \ ps)$

<proof>

lemma *length-mat-to-polys [simp]*: $\text{length } (\text{mat-to-polys } ts \ A) = \text{dim-row } A$

<proof>

lemma *mat-to-polys-nth*:

assumes $i < \text{dim-row } A$

shows $(\text{mat-to-polys } ts \ A) ! i = \text{row-to-poly } ts \ (\text{row } A \ i)$
 ⟨proof⟩

lemma *Keys-mat-to-polys*: $\text{Keys } (\text{set } (\text{mat-to-polys } ts \ A)) \subseteq \text{set } ts$
 ⟨proof⟩

lemma *polys-to-mat-to-polys*:
assumes $\text{Keys } (\text{set } ps) \subseteq \text{set } ts$
shows $\text{mat-to-polys } ts \ (\text{polys-to-mat } ts \ ps) = (ps::('t \Rightarrow_0 'b::\text{semiring-1}) \text{ list})$
 ⟨proof⟩

lemma *mat-to-polys-to-mat*:
assumes *distinct* ts **and** $\text{length } ts = \text{dim-col } A$
shows $(\text{polys-to-mat } ts \ (\text{mat-to-polys } ts \ A)) = A$
 ⟨proof⟩

15.4 Properties of Macaulay Matrices

lemma *row-to-poly-vec-times*:
assumes *distinct* ts **and** $\text{length } ts = \text{dim-col } A$
shows $\text{row-to-poly } ts \ (v \ v^* \ A) = ((\sum_{i=0..<\text{dim-row } A} (v \ \$ \ i) \cdot (\text{row-to-poly } ts \ (\text{row } A \ i))))::'t \Rightarrow_0 'b::\text{comm-semiring-0}$
 ⟨proof⟩

lemma *vec-times-polys-to-mat*:
assumes $\text{Keys } (\text{set } ps) \subseteq \text{set } ts$ **and** $v \in \text{carrier-vec } (\text{length } ps)$
shows $\text{row-to-poly } ts \ (v \ v^* \ (\text{polys-to-mat } ts \ ps)) = (\sum (c, p) \leftarrow \text{zip } (\text{list-of-vec } v) \ ps. \ c \cdot p)$
 (is $?l = ?r$)
 ⟨proof⟩

lemma *row-space-subset-phull*:
assumes $\text{Keys } (\text{set } ps) \subseteq \text{set } ts$
shows $\text{row-to-poly } ts \ ' \ \text{row-space } (\text{polys-to-mat } ts \ ps) \subseteq \text{phull } (\text{set } ps)$
 (is $?r \subseteq ?h$)
 ⟨proof⟩

lemma *phull-subset-row-space*:
assumes $\text{Keys } (\text{set } ps) \subseteq \text{set } ts$
shows $\text{phull } (\text{set } ps) \subseteq \text{row-to-poly } ts \ ' \ \text{row-space } (\text{polys-to-mat } ts \ ps)$
 (is $?h \subseteq ?r$)
 ⟨proof⟩

lemma *row-space-eq-phull*:
assumes $\text{Keys } (\text{set } ps) \subseteq \text{set } ts$
shows $\text{row-to-poly } ts \ ' \ \text{row-space } (\text{polys-to-mat } ts \ ps) = \text{phull } (\text{set } ps)$
 ⟨proof⟩

lemma *row-space-row-echelon-eq-phull*:

assumes $Keys (set\ ps) \subseteq set\ ts$
shows $row\text{-}to\text{-}poly\ ts \text{ ' row-space } (row\text{-}echelon (polys\text{-}to\text{-}mat\ ts\ ps)) = phull (set\ ps)$
 $\langle proof \rangle$

lemma *phull-row-echelon*:

assumes $Keys (set\ ps) \subseteq set\ ts$ **and** *distinct ts*
shows $phull (set (mat\text{-}to\text{-}polys\ ts (row\text{-}echelon (polys\text{-}to\text{-}mat\ ts\ ps)))) = phull (set\ ps)$
 $\langle proof \rangle$

lemma *pmdl-row-echelon*:

assumes $Keys (set\ ps) \subseteq set\ ts$ **and** *distinct ts*
shows $pmdl (set (mat\text{-}to\text{-}polys\ ts (row\text{-}echelon (polys\text{-}to\text{-}mat\ ts\ ps)))) = pmdl (set\ ps)$
(is ?l = ?r)
 $\langle proof \rangle$

end

context *ordered-term*

begin

lemma *lt-row-to-poly-pivot-fun*:

assumes $card\ S = dim\text{-}col (A::'b::semiring\text{-}1\ mat)$ **and** *pivot-fun A f (dim-col A)*
and $i < dim\text{-}row\ A$ **and** $f\ i < dim\text{-}col\ A$
shows $lt ((mat\text{-}to\text{-}polys (pps\text{-}to\text{-}list\ S)\ A) ! i) = (pps\text{-}to\text{-}list\ S) ! (f\ i)$
 $\langle proof \rangle$

lemma *lc-row-to-poly-pivot-fun*:

assumes $card\ S = dim\text{-}col (A::'b::semiring\text{-}1\ mat)$ **and** *pivot-fun A f (dim-col A)*
and $i < dim\text{-}row\ A$ **and** $f\ i < dim\text{-}col\ A$
shows $lc ((mat\text{-}to\text{-}polys (pps\text{-}to\text{-}list\ S)\ A) ! i) = 1$
 $\langle proof \rangle$

lemma *lt-row-to-poly-pivot-fun-less*:

assumes $card\ S = dim\text{-}col (A::'b::semiring\text{-}1\ mat)$ **and** *pivot-fun A f (dim-col A)*
and $i1 < i2$ **and** $i2 < dim\text{-}row\ A$ **and** $f\ i1 < dim\text{-}col\ A$ **and** $f\ i2 < dim\text{-}col\ A$
shows $(pps\text{-}to\text{-}list\ S) ! (f\ i2) \prec_t (pps\text{-}to\text{-}list\ S) ! (f\ i1)$
 $\langle proof \rangle$

lemma *lt-row-to-poly-pivot-fun-eqD*:

assumes $card\ S = dim\text{-}col (A::'b::semiring\text{-}1\ mat)$ **and** *pivot-fun A f (dim-col A)*
and $i1 < dim\text{-}row\ A$ **and** $i2 < dim\text{-}row\ A$ **and** $f\ i1 < dim\text{-}col\ A$ **and** $f\ i2 < dim\text{-}col\ A$

and $(pps\text{-to-list } S) ! (f\ i1) = (pps\text{-to-list } S) ! (f\ i2)$
shows $i1 = i2$
 $\langle proof \rangle$

lemma *lt-row-to-poly-pivot-in-keysD*:

assumes $card\ S = dim\text{-col } (A::'b::semiring-1\ mat)$ **and** $pivot\text{-fun } A\ f\ (dim\text{-col } A)$
and $i1 < dim\text{-row } A$ **and** $i2 < dim\text{-row } A$ **and** $f\ i1 < dim\text{-col } A$
and $(pps\text{-to-list } S) ! (f\ i1) \in keys\ ((mat\text{-to-polys } (pps\text{-to-list } S)\ A) ! i2)$
shows $i1 = i2$
 $\langle proof \rangle$

lemma *lt-row-space-pivot-fun*:

assumes $card\ S = dim\text{-col } (A::'b::\{comm\text{-semiring-0}, semiring-1\text{-no-zero-divisors}\} mat)$
and $pivot\text{-fun } A\ f\ (dim\text{-col } A)$ **and** $p \in row\text{-to-poly } (pps\text{-to-list } S) \text{ 'row-space } A$ **and** $p \neq 0$
shows $lt\ p \in lt\text{-set } (set\ (mat\text{-to-polys } (pps\text{-to-list } S)\ A))$
 $\langle proof \rangle$

15.5 Functions *Macaulay-mat* and *Macaulay-list*

definition *Macaulay-mat* :: $('t \Rightarrow_0 'b)\ list \Rightarrow 'b::field\ mat$

where $Macaulay\text{-mat } ps = polys\text{-to-mat } (Keys\text{-to-list } ps)\ ps$

definition *Macaulay-list* :: $('t \Rightarrow_0 'b)\ list \Rightarrow ('t \Rightarrow_0 'b::field)\ list$

where $Macaulay\text{-list } ps =$
 $filter\ (\lambda p. p \neq 0)\ (mat\text{-to-polys } (Keys\text{-to-list } ps)\ (row\text{-echelon } (Macaulay\text{-mat } ps)))$

lemma *dim-Macaulay-mat[simp]*:

$dim\text{-row } (Macaulay\text{-mat } ps) = length\ ps$

$dim\text{-col } (Macaulay\text{-mat } ps) = card\ (Keys\ (set\ ps))$

$\langle proof \rangle$

lemma *Macaulay-list-Nil [simp]*: $Macaulay\text{-list } [] = ([]::('t \Rightarrow_0 'b::field)\ list)$ (**is ?l = -**)

$\langle proof \rangle$

lemma *set-Macaulay-list*:

$set\ (Macaulay\text{-list } ps) =$

$set\ (mat\text{-to-polys } (Keys\text{-to-list } ps)\ (row\text{-echelon } (Macaulay\text{-mat } ps))) - \{0\}$

$\langle proof \rangle$

lemma *Keys-Macaulay-list*: $Keys\ (set\ (Macaulay\text{-list } ps)) \subseteq Keys\ (set\ ps)$

$\langle proof \rangle$

lemma *in-Macaulay-listE*:

assumes $p \in set\ (Macaulay\text{-list } ps)$

and *pivot-fun* (row-echelon (Macaulay-mat ps)) f (dim-col (row-echelon (Macaulay-mat ps)))
obtains *i* **where** $i < \text{dim-row (row-echelon (Macaulay-mat ps))}$
and $p = (\text{mat-to-polys (Keys-to-list ps) (row-echelon (Macaulay-mat ps))}) ! i$
and $f i < \text{dim-col (row-echelon (Macaulay-mat ps))}$
 <proof>

lemma *phull-Macaulay-list*: $\text{phull (set (Macaulay-list ps))} = \text{phull (set ps)}$
 <proof>

lemma *pmdl-Macaulay-list*: $\text{pmdl (set (Macaulay-list ps))} = \text{pmdl (set ps)}$
 <proof>

lemma *Macaulay-list-is-monic-set*: $\text{is-monic-set (set (Macaulay-list ps))}$
 <proof>

lemma *Macaulay-list-not-zero*: $0 \notin \text{set (Macaulay-list ps)}$
 <proof>

lemma *Macaulay-list-distinct-lt*:
assumes $x \in \text{set (Macaulay-list ps)}$ **and** $y \in \text{set (Macaulay-list ps)}$
and $x \neq y$
shows $\text{lt } x \neq \text{lt } y$
 <proof>

lemma *Macaulay-list-lt*:
assumes $p \in \text{phull (set ps)}$ **and** $p \neq 0$
obtains g **where** $g \in \text{set (Macaulay-list ps)}$ **and** $g \neq 0$ **and** $\text{lt } p = \text{lt } g$
 <proof>

end

end

16 Faugère's F4 Algorithm

theory *F4*
imports *Macaulay-Matrix Algorithm-Schema*
begin

This theory implements Faugère's F4 algorithm based on *gd-term.gb-schema-direct*.

16.1 Symbolic Preprocessing

context *gd-term*
begin

definition *sym-preproc-aux-term1* :: $('a \Rightarrow \text{nat}) \Rightarrow (((t \Rightarrow_0 'b) \text{ list} \times 't \text{ list} \times 't \text{ list} \times (t \Rightarrow_0 'b) \text{ list}) \times$

($'t \Rightarrow_0 'b$) list \times $'t$ list \times $'t$ list \times ($'t \Rightarrow_0$
 $'b$) list)) set
where *sym-preproc-aux-term1* $d =$
 $\{((gs1, ks1, ts1, fs1), (gs2::('t \Rightarrow_0 'b) list, ks2, ts2, fs2)). \exists t2 \in set\ ts2.$
 $\forall t1 \in set\ ts1. t1 \prec_t t2\}$

definition *sym-preproc-aux-term2* $:: ('a \Rightarrow nat) \Rightarrow ((('t \Rightarrow_0 'b::zero) list \times 't list$
 $\times 't list \times ('t \Rightarrow_0 'b) list) \times$
 $(('t \Rightarrow_0 'b) list \times 't list \times 't list \times ('t \Rightarrow_0$
 $'b) list)) set$
where *sym-preproc-aux-term2* $d =$
 $\{((gs1, ks1, ts1, fs1), (gs2::('t \Rightarrow_0 'b) list, ks2, ts2, fs2)). gs1 = gs2 \wedge$
 $dgrad-set-le\ d\ (pp-of-term\ ' set\ ts1)\ (pp-of-term$
 $' (Keys\ (set\ gs2) \cup set\ ts2))\}$

definition *sym-preproc-aux-term*
where *sym-preproc-aux-term* $d = sym-preproc-aux-term1\ d \cap sym-preproc-aux-term2$
 d

lemma *wfp-on-ord-term-strict*:
assumes *dickson-grading* d
shows *wfp-on* (\prec_t) (*pp-of-term* - $' dgrad-set\ d\ m$)
 $\langle proof \rangle$

lemma *sym-preproc-aux-term1-wf-on*:
assumes *dickson-grading* d
shows *wfp-on* ($\lambda x\ y. (x, y) \in sym-preproc-aux-term1\ d$) $\{x. set\ (fst\ (snd\ (snd$
 $x))) \subseteq pp-of-term\ - ' dgrad-set\ d\ m\}$
 $\langle proof \rangle$

lemma *sym-preproc-aux-term-wf*:
assumes *dickson-grading* d
shows *wf* (*sym-preproc-aux-term* d)
 $\langle proof \rangle$

primrec *sym-preproc-addnew* $:: ('t \Rightarrow_0 'b::semiring-1) list \Rightarrow 't list \Rightarrow ('t \Rightarrow_0 'b)$
 $list \Rightarrow 't \Rightarrow$

($'t list \times ('t \Rightarrow_0 'b) list$) **where**
 $sym-preproc-addnew\ []\ vs\ fs\ - = (vs, fs)|$
 $sym-preproc-addnew\ (g\ \# gs)\ vs\ fs\ v =$
 $(if\ lt\ g\ adds_t\ v\ then$
 $(let\ f = monom-mult\ 1\ (pp-of-term\ v - lp\ g)\ g\ in$
 $sym-preproc-addnew\ gs\ (merge-wrt\ (\succ_t)\ vs\ (keys-to-list\ (tail\ f)))\ (insert-list$
 $f\ fs)\ v$
 $)$
 $else$
 $sym-preproc-addnew\ gs\ vs\ fs\ v$
 $)$

lemma *fst-sym-preproc-addnew-less:*

assumes $\bigwedge u. u \in \text{set } vs \implies u \prec_t v$
and $u \in \text{set } (\text{fst } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$
shows $u \prec_t v$
<proof>

lemma *fst-sym-preproc-addnew-dgrad-set-le:*

assumes *dickson-grading* d
shows *dgrad-set-le* d (*pp-of-term* ‘ $\text{set } (\text{fst } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$ ’)
(*pp-of-term* ‘ $(\text{Keys } (\text{set } gs) \cup \text{insert } v \ (\text{set } vs))$ ’)
<proof>

lemma *components-fst-sym-preproc-addnew-subset:*

component-of-term ‘ $\text{set } (\text{fst } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$ ’ \subseteq *component-of-term*
‘ $(\text{Keys } (\text{set } gs) \cup \text{insert } v \ (\text{set } vs))$ ’
<proof>

lemma *fst-sym-preproc-addnew-superset:* $\text{set } vs \subseteq \text{set } (\text{fst } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$
<proof>

lemma *snd-sym-preproc-addnew-superset:* $\text{set } fs \subseteq \text{set } (\text{snd } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$
<proof>

lemma *in-snd-sym-preproc-addnewE:*

assumes $p \in \text{set } (\text{snd } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$
assumes *1:* $p \in \text{set } fs \implies \text{thesis}$
assumes *2:* $\bigwedge g \ s. g \in \text{set } gs \implies p = \text{monom-mult } 1 \ s \ g \implies \text{thesis}$
shows *thesis*
<proof>

lemma *sym-preproc-addnew-pmdl:*

$\text{pmdl } (\text{set } gs \cup \text{set } (\text{snd } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))) = \text{pmdl } (\text{set } gs \cup \text{set } fs)$
(**is** $\text{pmdl } (\text{set } gs \cup ?l) = ?r$)
<proof>

lemma *Keys-snd-sym-preproc-addnew:*

$\text{Keys } (\text{set } (\text{snd } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))) \cup \text{insert } v \ (\text{set } vs) =$
 $\text{Keys } (\text{set } fs) \cup \text{insert } v \ (\text{set } (\text{fst } (\text{sym-preproc-addnew } gs \text{ vs } (fs::('t \Rightarrow_0 'b::\text{semiring-1-no-zero-divisors})$
list) $v)))$
<proof>

lemma *sym-preproc-addnew-complete:*

assumes $g \in \text{set } gs$ **and** $lt \ g \ \text{adds}_t \ v$
shows $\text{monom-mult } 1 \ (\text{pp-of-term } v - lp \ g) \ g \in \text{set } (\text{snd } (\text{sym-preproc-addnew } gs \text{ vs } fs \ v))$
<proof>

function *sym-preproc-aux* :: ('t \Rightarrow_0 'b::semiring-1) list \Rightarrow 't list \Rightarrow ('t list \times ('t \Rightarrow_0 'b) list) \Rightarrow
('t list \times ('t \Rightarrow_0 'b) list) **where**
sym-preproc-aux gs ks (vs, fs) =
(if vs = [] then
(ks, fs)
else
let v = ord-term-lin.max-list vs; vs' = removeAll v vs in
sym-preproc-aux gs (ks @ [v]) (*sym-preproc-addnew* gs vs' fs v)
)
⟨proof⟩
termination ⟨proof⟩

lemma *sym-preproc-aux-Nil*: *sym-preproc-aux* gs ks ([], fs) = (ks, fs)
⟨proof⟩

lemma *sym-preproc-aux-sorted*:
assumes sorted-wrt (\succ_t) (v # vs)
shows *sym-preproc-aux* gs ks (v # vs, fs) = *sym-preproc-aux* gs (ks @ [v])
(*sym-preproc-addnew* gs vs fs v)
⟨proof⟩

lemma *sym-preproc-aux-induct* [consumes 0, case-names base rec]:
assumes base: \bigwedge ks fs. P ks [] fs (ks, fs)
and rec: \bigwedge ks vs fs v vs'. vs \neq [] \implies v = ord-term-lin.Max (set vs) \implies vs' =
removeAll v vs \implies
P (ks @ [v]) (fst (*sym-preproc-addnew* gs vs' fs v)) (snd (*sym-preproc-addnew*
gs vs' fs v))
(*sym-preproc-aux* gs (ks @ [v]) (*sym-preproc-addnew* gs vs' fs v))
 \implies
P ks vs fs (*sym-preproc-aux* gs (ks @ [v]) (*sym-preproc-addnew* gs vs'
fs v))
shows P ks vs fs (*sym-preproc-aux* gs ks (vs, fs))
⟨proof⟩

lemma *fst-sym-preproc-aux-sorted-wrt*:
assumes sorted-wrt (\succ_t) ks **and** \bigwedge k v. k \in set ks \implies v \in set vs \implies v \prec_t k
shows sorted-wrt (\succ_t) (fst (*sym-preproc-aux* gs ks (vs, fs)))
⟨proof⟩

lemma *fst-sym-preproc-aux-complete*:
assumes Keys (set (fs::('t \Rightarrow_0 'b::semiring-1-no-zero-divisors) list)) = set ks \cup
set vs
shows set (fst (*sym-preproc-aux* gs ks (vs, fs))) = Keys (set (snd (*sym-preproc-aux*
gs ks (vs, fs))))
⟨proof⟩

lemma *snd-sym-preproc-aux-superset*: set fs \subseteq set (snd (*sym-preproc-aux* gs ks (vs,

$fs)))$
 $\langle proof \rangle$

lemma *in-snd-sym-preproc-auxE*:

assumes $p \in \text{set } (snd \ (sym\text{-preproc-aux } gs \ ks \ (vs, fs)))$

assumes 1: $p \in \text{set } fs \implies \text{thesis}$

assumes 2: $\bigwedge g \ t. g \in \text{set } gs \implies p = \text{monom-mult } 1 \ t \ g \implies \text{thesis}$

shows *thesis*

$\langle proof \rangle$

lemma *snd-sym-preproc-aux-pmdl*:

$pmdl \ (\text{set } gs \cup \text{set } (snd \ (sym\text{-preproc-aux } gs \ ks \ (ts, fs)))) = pmdl \ (\text{set } gs \cup \text{set } fs)$

$\langle proof \rangle$

lemma *snd-sym-preproc-aux-dgrad-set-le*:

assumes *dickson-grading* d **and** $\text{set } vs \subseteq \text{Keys} \ (\text{set } (fs::('t \Rightarrow_0 'b::\text{semiring-1-no-zero-divisors}) \text{ list}))$

shows *dgrad-set-le* d $(pp\text{-of-term } ' \text{Keys} \ (\text{set } (snd \ (sym\text{-preproc-aux } gs \ ks \ (vs, fs)))))) \ (pp\text{-of-term } ' \text{Keys} \ (\text{set } gs \cup \text{set } fs))$

$\langle proof \rangle$

lemma *components-snd-sym-preproc-aux-subset*:

assumes $\text{set } vs \subseteq \text{Keys} \ (\text{set } (fs::('t \Rightarrow_0 'b::\text{semiring-1-no-zero-divisors}) \text{ list}))$

shows *component-of-term* $' \text{Keys} \ (\text{set } (snd \ (sym\text{-preproc-aux } gs \ ks \ (vs, fs)))) \subseteq$
component-of-term $' \text{Keys} \ (\text{set } gs \cup \text{set } fs)$

$\langle proof \rangle$

lemma *snd-sym-preproc-aux-complete*:

assumes $\bigwedge u' \ g'. u' \in \text{Keys} \ (\text{set } fs) \implies u' \notin \text{set } vs \implies g' \in \text{set } gs \implies lt \ g' \text{ adds}_t \ u' \implies$

$\text{monom-mult } 1 \ (pp\text{-of-term } u' - lp \ g') \ g' \in \text{set } fs$

assumes $u \in \text{Keys} \ (\text{set } (snd \ (sym\text{-preproc-aux } gs \ ks \ (vs, fs))))$ **and** $g \in \text{set } gs$
and $lt \ g \text{ adds}_t \ u$

shows $\text{monom-mult} \ (1::'b::\text{semiring-1-no-zero-divisors}) \ (pp\text{-of-term } u - lp \ g) \ g \in$

$\text{set } (snd \ (sym\text{-preproc-aux } gs \ ks \ (vs, fs)))$

$\langle proof \rangle$

definition *sym-preproc* $:: ('t \Rightarrow_0 'b::\text{semiring-1}) \text{ list} \Rightarrow ('t \Rightarrow_0 'b) \text{ list} \Rightarrow ('t \text{ list} \times ('t \Rightarrow_0 'b) \text{ list})$

where $sym\text{-preproc } gs \ fs = sym\text{-preproc-aux } gs \ [] \ (\text{Keys-to-list } fs, fs)$

lemma *sym-preproc-Nil [simp]*: $sym\text{-preproc } gs \ [] = ([], [])$

$\langle proof \rangle$

lemma *fst-sym-preproc*:

$fst \ (sym\text{-preproc } gs \ fs) = \text{Keys-to-list} \ (snd \ (sym\text{-preproc } gs \ (fs::('t \Rightarrow_0 'b::\text{semiring-1-no-zero-divisors}) \text{ list})))$

<proof>

lemma *snd-sym-preproc-superset*: $set\ fs \subseteq set\ (snd\ (sym\ preproc\ gs\ fs))$
<proof>

lemma *in-snd-sym-preprocE*:

assumes $p \in set\ (snd\ (sym\ preproc\ gs\ fs))$

assumes 1: $p \in set\ fs \implies thesis$

assumes 2: $\bigwedge g\ t.\ g \in set\ gs \implies p = monom\ mult\ 1\ t\ g \implies thesis$

shows *thesis*

<proof>

lemma *snd-sym-preproc-pmdl*: $pmdl\ (set\ gs \cup set\ (snd\ (sym\ preproc\ gs\ fs))) =$
 $pmdl\ (set\ gs \cup set\ fs)$

<proof>

lemma *snd-sym-preproc-dgrad-set-le*:

assumes *dickson-grading* d

shows $dgrad\ set\ le\ d\ (pp\ of\ term\ 'Keys\ (set\ (snd\ (sym\ preproc\ gs\ fs))))$

$(pp\ of\ term\ 'Keys\ (set\ gs \cup set\ (fs::('t \Rightarrow_0\ 'b::semiring-1-no-zero-divisors)$

$list)))$

<proof>

corollary *snd-sym-preproc-dgrad-p-set-le*:

assumes *dickson-grading* d

shows $dgrad\ p\ set\ le\ d\ (set\ (snd\ (sym\ preproc\ gs\ fs)))\ (set\ gs \cup set\ (fs::('t \Rightarrow_0$
 $'b::semiring-1-no-zero-divisors)\ list))$

<proof>

lemma *components-snd-sym-preproc-subset*:

$component\ of\ term\ 'Keys\ (set\ (snd\ (sym\ preproc\ gs\ fs))) \subseteq$

$component\ of\ term\ 'Keys\ (set\ gs \cup set\ (fs::('t \Rightarrow_0\ 'b::semiring-1-no-zero-divisors)$

$list))$

<proof>

lemma *snd-sym-preproc-complete*:

assumes $v \in Keys\ (set\ (snd\ (sym\ preproc\ gs\ fs)))$ **and** $g \in set\ gs$ **and** $lt\ g\ adds_t$
 v

shows $monom\ mult\ (1::'b::semiring-1-no-zero-divisors)\ (pp\ of\ term\ v - lp\ g)\ g$
 $\in set\ (snd\ (sym\ preproc\ gs\ fs))$

<proof>

end

16.2 *lin-red*

context *ordered-term*

begin

definition $lin-red :: ('t \Rightarrow_0 'b::field) set \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow ('t \Rightarrow_0 'b) \Rightarrow bool$
where $lin-red F p q \equiv (\exists f \in F. red-single p q f 0)$

$lin-red$ is a restriction of red , where the reductor (f) may only be multiplied by a constant factor, i. e. where the power-product is 0 .

lemma $lin-redI$:
assumes $f \in F$ **and** $red-single p q f 0$
shows $lin-red F p q$
 $\langle proof \rangle$

lemma $lin-redE$:
assumes $lin-red F p q$
obtains $f::'t \Rightarrow_0 'b::field$ **where** $f \in F$ **and** $red-single p q f 0$
 $\langle proof \rangle$

lemma $lin-red-imp-red$:
assumes $lin-red F p q$
shows $red F p q$
 $\langle proof \rangle$

lemma $lin-red-Un$: $lin-red (F \cup G) p q = (lin-red F p q \vee lin-red G p q)$
 $\langle proof \rangle$

lemma $lin-red-imp-red-rtrancl$:
assumes $(lin-red F)^{**} p q$
shows $(red F)^{**} p q$
 $\langle proof \rangle$

lemma $phull-closed-lin-red$:
assumes $phull B \subseteq phull A$ **and** $p \in phull A$ **and** $lin-red B p q$
shows $q \in phull A$
 $\langle proof \rangle$

16.3 Reduction

definition $Macaulay-red :: 't list \Rightarrow ('t \Rightarrow_0 'b) list \Rightarrow ('t \Rightarrow_0 'b::field) list$
where $Macaulay-red vs fs =$
 $(let lts = map lt (filter (\lambda p. p \neq 0) fs) in$
 $filter (\lambda p. p \neq 0 \wedge lt p \notin set lts) (mat-to-polys vs (row-echelon (polys-to-mat$
 $vs fs))))$
 $)$

$Macaulay-red vs fs$ auto-reduces (w. r. t. $lin-red$) the given list fs and returns those non-zero polynomials whose leading terms are not in $lt-set$ ($set fs$). Argument vs is expected to be $Keys-to-list fs$; this list is passed as an argument to $Macaulay-red$, because it can be efficiently computed by symbolic preprocessing.

lemma $Macaulay-red-alt$:

$\text{Macaulay-red (Keys-to-list fs) fs} = \text{filter } (\lambda p. \text{lt } p \notin \text{lt-set (set fs)}) (\text{Macaulay-list fs})$
 ⟨proof⟩

lemma *set-Macaulay-red*:

$\text{set (Macaulay-red (Keys-to-list fs) fs)} = \text{set (Macaulay-list fs)} - \{p. \text{lt } p \in \text{lt-set (set fs)}\}$
 ⟨proof⟩

lemma *Keys-Macaulay-red*: $\text{Keys (set (Macaulay-red (Keys-to-list fs) fs))} \subseteq \text{Keys (set fs)}$
 ⟨proof⟩

end

context *gd-term*

begin

lemma *Macaulay-red-reducible*:

assumes $f \in \text{phull (set fs)}$ **and** $F \subseteq \text{set fs}$ **and** $\text{lt-set } F = \text{lt-set (set fs)}$
shows $(\text{lin-red (F } \cup \text{ set (Macaulay-red (Keys-to-list fs) fs))})^{**} f 0$
 ⟨proof⟩

primrec *pdata-pairs-to-list* :: $(t, 'b::\text{field}, 'c) \text{pdata-pair list} \Rightarrow (t \Rightarrow_0 'b) \text{list}$
where

$\text{pdata-pairs-to-list } [] = []$
 $\text{pdata-pairs-to-list (p \# ps)} =$
 (let $f = \text{fst (fst p)}$; $g = \text{fst (snd p)}$; $lf = \text{lp } f$; $lg = \text{lp } g$; $l = \text{lcs } lf \text{ } lg$ in
 (monom-mult (1 / lc f) (l - lf) f) # (monom-mult (1 / lc g) (l - lg) g) #
 (pdata-pairs-to-list ps)
)

lemma *in-pdata-pairs-to-listI1*:

assumes $(f, g) \in \text{set ps}$
shows $\text{monom-mult (1 / lc (fst f)) ((lcs (lp (fst f)) (lp (fst g))) - (lp (fst f)))}$
 $(fst f) \in \text{set (pdata-pairs-to-list ps)}$ **(is ?m ∈ -)**
 ⟨proof⟩

lemma *in-pdata-pairs-to-listI2*:

assumes $(f, g) \in \text{set ps}$
shows $\text{monom-mult (1 / lc (fst g)) ((lcs (lp (fst f)) (lp (fst g))) - (lp (fst g)))}$
 $(fst g) \in \text{set (pdata-pairs-to-list ps)}$ **(is ?m ∈ -)**
 ⟨proof⟩

lemma *in-pdata-pairs-to-listE*:

assumes $h \in \text{set (pdata-pairs-to-list ps)}$
obtains $f \text{ } g$ **where** $(f, g) \in \text{set ps} \vee (g, f) \in \text{set ps}$
and $h = \text{monom-mult (1 / lc (fst f)) ((lcs (lp (fst f)) (lp (fst g))) - (lp (fst f)))}$
 $(fst f)$

<proof>

definition $f_4\text{-red-}aux :: ('t, 'b::field, 'c) \text{pdata list} \Rightarrow ('t, 'b, 'c) \text{pdata-pair list} \Rightarrow ('t \Rightarrow_0 'b) \text{list}$

where $f_4\text{-red-}aux \text{ bs ps} =$
 $(\text{let } aux = \text{sym-preproc } (\text{map } fst \text{ bs}) (\text{pdata-pairs-to-list } ps) \text{ in Macaulay-red } (fst \text{ aux}) (\text{snd } aux))$

$f_4\text{-red-}aux$ only takes two arguments, since it does not distinguish between those elements of the current basis that are known to be a Gröbner basis (called gs in *Groebner-Bases.Algorithm-Schema*) and the remaining ones.

lemma $f_4\text{-red-}aux\text{-not-zero}: 0 \notin \text{set } (f_4\text{-red-}aux \text{ bs ps})$

<proof>

lemma $f_4\text{-red-}aux\text{-irreducible}$:

assumes $h \in \text{set } (f_4\text{-red-}aux \text{ bs ps})$ **and** $b \in \text{set } bs$ **and** $fst \ b \neq 0$

shows $\neg lt \ (fst \ b) \ \text{adds}_t \ lt \ h$

<proof>

lemma $f_4\text{-red-}aux\text{-dgrad-p-set-le}$:

assumes $dickson\text{-grading } d$

shows $dgrad\text{-p-set-le } d \ (\text{set } (f_4\text{-red-}aux \text{ bs ps})) \ (\text{args-to-set } ([], bs, ps))$

<proof>

lemma $components\text{-}f_4\text{-red-}aux\text{-subset}$:

$component\text{-of-term } ' \ Keys \ (\text{set } (f_4\text{-red-}aux \text{ bs ps})) \subseteq component\text{-of-term } ' \ Keys \ (\text{args-to-set } ([], bs, ps))$

<proof>

lemma $pmdl\text{-}f_4\text{-red-}aux$: $\text{set } (f_4\text{-red-}aux \text{ bs ps}) \subseteq pmdl \ (\text{args-to-set } ([], bs, ps))$

<proof>

lemma $f_4\text{-red-}aux\text{-phull-reducible}$:

assumes $\text{set } ps \subseteq \text{set } bs \times \text{set } bs$

and $f \in phull \ (\text{set } (\text{pdata-pairs-to-list } ps))$

shows $(\text{red } (fst \ ' \ \text{set } bs \cup \text{set } (f_4\text{-red-}aux \text{ bs ps})))^{**} f \ 0$

<proof>

corollary $f_4\text{-red-}aux\text{-spoly-reducible}$:

assumes $\text{set } ps \subseteq \text{set } bs \times \text{set } bs$ **and** $(p, q) \in \text{set } ps$

shows $(\text{red } (fst \ ' \ \text{set } bs \cup \text{set } (f_4\text{-red-}aux \text{ bs ps})))^{**} (\text{spoly } (fst \ p) \ (fst \ q)) \ 0$

<proof>

definition $f_4\text{-red} :: ('t, 'b::field, 'c::default, 'd) \text{compl}T$

where $f_4\text{-red } gs \text{ bs } ps \text{ sps } data = (\text{map } (\lambda h. (h, default)) (f_4\text{-red-}aux \ (gs \ @ \ bs) \ sps), \text{snd } data)$

lemma $fst\text{-set-}fst\text{-}f_4\text{-red}$: $fst \ ' \ \text{set } (fst \ (f_4\text{-red } gs \ bs \ ps \ sps \ data)) = \text{set } (f_4\text{-red-}aux \ (gs \ @ \ bs) \ sps)$

<proof>

lemma *rcp-spec-f4-red*: *rcp-spec f4-red*
<proof>

lemmas *compl-struct-f4-red* = *compl-struct-rcp[OF rcp-spec-f4-red]*

lemmas *compl-pmdl-f4-red* = *compl-pmdl-rcp[OF rcp-spec-f4-red]*

lemmas *compl-conn-f4-red* = *compl-conn-rcp[OF rcp-spec-f4-red]*

16.4 Pair Selection

primrec *f4-sel-aux* :: 'a \Rightarrow ('t, 'b::zero, 'c) *pdata-pair list* \Rightarrow ('t, 'b, 'c) *pdata-pair list* **where**
f4-sel-aux - [] = []
f4-sel-aux t (p # ps) =
 (if (lcs (lp (fst (fst p))) (lp (fst (snd p)))) = t then
 p # (*f4-sel-aux* t ps)
 else
 []
)

lemma *f4-sel-aux-subset*: *set (f4-sel-aux t ps) \subseteq set ps*
<proof>

primrec *f4-sel* :: ('t, 'b::zero, 'c, 'd) *selT* **where**
f4-sel gs bs [] data = []
f4-sel gs bs (p # ps) data = p # (*f4-sel-aux* (lcs (lp (fst (fst p))) (lp (fst (snd p)))) ps)

lemma *sel-spec-f4-sel*: *sel-spec f4-sel*
<proof>

16.5 The F4 Algorithm

The F4 algorithm is just *gb-schema-direct* with parameters instantiated by suitable functions.

lemma *struct-spec-f4*: *struct-spec f4-sel add-pairs-canon add-basis-canon f4-red*
<proof>

definition *f4-aux* :: ('t, 'b, 'c) *pdata list* \Rightarrow $\text{nat} \times \text{nat} \times 'd \Rightarrow$ ('t, 'b, 'c) *pdata list*
 \Rightarrow

('t, 'b, 'c) *pdata-pair list* \Rightarrow ('t, 'b::field, 'c::default) *pdata list*
 where *f4-aux* = *gb-schema-aux f4-sel add-pairs-canon add-basis-canon f4-red*

lemmas *f4-aux-simps* [code] = *gb-schema-aux-simps[OF struct-spec-f4, folded f4-aux-def]*

definition *f4* :: ('t, 'b, 'c) *pdata' list* \Rightarrow 'd \Rightarrow ('t, 'b::field, 'c::default) *pdata' list*
 where *f4* = *gb-schema-direct f4-sel add-pairs-canon add-basis-canon f4-red*

lemmas *f4-simps* [code] = *gb-schema-direct-def*[of *f4-sel add-pairs-canon add-basis-canon f4-red*, folded *f4-def f4-aux-def*]

lemmas *f4-isGB* = *gb-schema-direct-isGB*[OF *struct-spec-f4 compl-conn-f4-red*, folded *f4-def*]

lemmas *f4-pmdl* = *gb-schema-direct-pmdl*[OF *struct-spec-f4 compl-pmdl-f4-red*, folded *f4-def*]

16.5.1 Special Case: *punit*

lemma (in *gd-term*) *struct-spec-f4-punit*: *punit.struct-spec punit.f4-sel add-pairs-punit-canon punit.add-basis-canon punit.f4-red*
 ⟨*proof*⟩

definition *f4-aux-punit* :: ('a, 'b, 'c) *pdata list* ⇒ *nat* × *nat* × 'd ⇒ ('a, 'b, 'c) *pdata list* ⇒
 ('a, 'b, 'c) *pdata-pair list* ⇒ ('a, 'b::field, 'c::default) *pdata list*

where *f4-aux-punit* = *punit.gb-schema-aux punit.f4-sel add-pairs-punit-canon punit.add-basis-canon punit.f4-red*

lemmas *f4-aux-punit-simps* [code] = *punit.gb-schema-aux-simps*[OF *struct-spec-f4-punit*, folded *f4-aux-punit-def*]

definition *f4-punit* :: ('a, 'b, 'c) *pdata' list* ⇒ 'd ⇒ ('a, 'b::field, 'c::default) *pdata' list*

where *f4-punit* = *punit.gb-schema-direct punit.f4-sel add-pairs-punit-canon punit.add-basis-canon punit.f4-red*

lemmas *f4-punit-simps* [code] = *punit.gb-schema-direct-def*[of *punit.f4-sel add-pairs-punit-canon punit.add-basis-canon punit.f4-red*, folded *f4-punit-def f4-aux-punit-def*]

lemmas *f4-punit-isGB* = *punit.gb-schema-direct-isGB*[OF *struct-spec-f4-punit punit.compl-conn-f4-red*, folded *f4-punit-def*]

lemmas *f4-punit-pmdl* = *punit.gb-schema-direct-pmdl*[OF *struct-spec-f4-punit punit.compl-pmdl-f4-red*, folded *f4-punit-def*]

end

end

17 Sample Computations with the F4 Algorithm

theory *F4-Examples*

imports *F4 Algorithm-Schema-Impl Jordan-Normal-Form. Gauss-Jordan-IArray-Impl Code-Target-Rat*

begin

We only consider scalar polynomials here, but vector-polynomials could be handled, too.

17.1 Preparations

primrec *remdups-wrt-rev* :: ('a ⇒ 'b) ⇒ 'a list ⇒ 'b list ⇒ 'a list **where**
remdups-wrt-rev f [] vs = [] |
remdups-wrt-rev f (x # xs) vs =
 (let fx = f x in if List.member vs fx then *remdups-wrt-rev* f xs vs else x #
 (*remdups-wrt-rev* f xs (fx # vs)))

lemma *remdups-wrt-rev-notin*: $v \in \text{set } vs \implies v \notin f' \text{ set } (\text{remdups-wrt-rev } f \text{ xs } vs)$
 <proof>

lemma *distinct-remdups-wrt-rev*: *distinct* (map f (*remdups-wrt-rev* f xs vs))
 <proof>

lemma *map-of-remdups-wrt-rev'*:
 map-of (*remdups-wrt-rev* fst xs vs) k = map-of (filter ($\lambda x. \text{fst } x \notin \text{set } vs$) xs) k
 <proof>

corollary *map-of-remdups-wrt-rev*: map-of (*remdups-wrt-rev* fst xs []) = map-of
 xs
 <proof>

lemma (in *term-powerprod*) *compute-list-to-poly* [code]:
list-to-poly ts cs = distr₀ DRLEX (*remdups-wrt-rev* fst (zip ts cs) [])
 <proof>

lemma (in *ordered-term*) *compute-Macaulay-list* [code]:
Macaulay-list ps =
 (let ts = *Keys-to-list* ps in
 filter ($\lambda p. p \neq 0$) (mat-to-polys ts (row-echelon (polys-to-mat ts ps))))
)
 <proof>

declare *conversep-iff* [code]

derive (eq) *ceq poly-mapping*
derive (no) *ccompare poly-mapping*
derive (dlist) *set-impl poly-mapping*
derive (no) *cenum poly-mapping*

derive (eq) *ceq rat*
derive (no) *ccompare rat*
derive (dlist) *set-impl rat*
derive (no) *cenum rat*

global-interpretation *punit'*: *gd-powerprod ord-pp-punit cmp-term ord-pp-strict-punit*

cmp-term

```
rewrites punit.adds-term = (adds)
and punit.pp-of-term = ( $\lambda x. x$ )
and punit.component-of-term = ( $\lambda-. ()$ )
and punit.monom-mult = monom-mult-punit
and punit.mult-scalar = mult-scalar-punit
and punit'.punit.min-term = min-term-punit
and punit'.punit.lt = lt-punit cmp-term
and punit'.punit.lc = lc-punit cmp-term
and punit'.punit.tail = tail-punit cmp-term
and punit'.punit.ord-p = ord-p-punit cmp-term
and punit'.punit.ord-strict-p = ord-strict-p-punit cmp-term
and punit'.punit.keys-to-list = keys-to-list-punit cmp-term
for cmp-term :: ('a::nat, 'b::{nat,add-wellorder}) pp nat-term-order
```

```
defines max-punit = punit'.ordered-powerprod-lin.max
and max-list-punit = punit'.ordered-powerprod-lin.max-list
and find-adds-punit = punit'.punit.find-adds
and trd-aux-punit = punit'.punit.trd-aux
and trd-punit = punit'.punit.trd
and spoly-punit = punit'.punit.spoly
and count-const-lt-components-punit = punit'.punit.count-const-lt-components
and count-rem-components-punit = punit'.punit.count-rem-components
and const-lt-component-punit = punit'.punit.const-lt-component
and full-gb-punit = punit'.punit.full-gb
and add-pairs-single-sorted-punit = punit'.punit.add-pairs-single-sorted
and add-pairs-punit = punit'.punit.add-pairs
and canon-pair-order-aux-punit = punit'.punit.canon-pair-order-aux
and canon-basis-order-punit = punit'.punit.canon-basis-order
and new-pairs-sorted-punit = punit'.punit.new-pairs-sorted
and product-crit-punit = punit'.punit.product-crit
and chain-ncrit-punit = punit'.punit.chain-ncrit
and chain-ocrit-punit = punit'.punit.chain-ocrit
and apply-icrit-punit = punit'.punit.apply-icrit
and apply-ncrit-punit = punit'.punit.apply-ncrit
and apply-ocrit-punit = punit'.punit.apply-ocrit
and Keys-to-list-punit = punit'.punit.Keys-to-list
and sym-preproc-addnew-punit = punit'.punit.sym-preproc-addnew
and sym-preproc-aux-punit = punit'.punit.sym-preproc-aux
and sym-preproc-punit = punit'.punit.sym-preproc
and Macaulay-mat-punit = punit'.punit.Macaulay-mat
and Macaulay-list-punit = punit'.punit.Macaulay-list
and pdata-pairs-to-list-punit = punit'.punit.pdata-pairs-to-list
and Macaulay-red-punit = punit'.punit.Macaulay-red
and f4-sel-aux-punit = punit'.punit.f4-sel-aux
and f4-sel-punit = punit'.punit.f4-sel
and f4-red-aux-punit = punit'.punit.f4-red-aux
and f4-red-punit = punit'.punit.f4-red
and f4-aux-punit = punit'.punit.f4-aux-punit
```

and $f_4\text{-punit} = \text{punit}'.\text{punit}.f_4\text{-punit}$
 $\langle \text{proof} \rangle$

17.2 Computations

experiment begin interpretation $\text{trivariate}_0\text{-rat}$ $\langle \text{proof} \rangle$

lemma

$lt\text{-punit DRLEX } (X^2 * Z^3 + 3 * X^2 * Y) = \text{sparse}_0 [(0, 2), (2, 3)]$
 $\langle \text{proof} \rangle$

lemma

$lc\text{-punit DRLEX } (X^2 * Z^3 + 3 * X^2 * Y) = 1$
 $\langle \text{proof} \rangle$

lemma

$tail\text{-punit DRLEX } (X^2 * Z^3 + 3 * X^2 * Y) = 3 * X^2 * Y$
 $\langle \text{proof} \rangle$

lemma

$ord\text{-strict-p-punit DRLEX } (X^2 * Z^4 - 2 * Y^3 * Z^2) (X^2 * Z^7 + 2 * Y^3 * Z^2)$
 $\langle \text{proof} \rangle$

lemma

$f_4\text{-punit DRLEX}$
 $[$
 $(X^2 * Z^4 - 2 * Y^3 * Z^2, ()),$
 $(Y^2 * Z + 2 * Z^3, ()),$
 $]$ $() =$
 $[$
 $(X^2 * Y^2 * Z^2 + 4 * Y^3 * Z^2, ()),$
 $(X^2 * Z^4 - 2 * Y^3 * Z^2, ()),$
 $(Y^2 * Z + 2 * Z^3, ()),$
 $(X^2 * Y^4 * Z + 4 * Y^5 * Z, ()),$
 $]$
 $\langle \text{proof} \rangle$

lemma

$f_4\text{-punit DRLEX}$
 $[$
 $(X^2 + Y^2 + Z^2 - 1, ()),$
 $(X * Y - Z - 1, ()),$
 $(Y^2 + X, ()),$
 $(Z^2 + X, ()),$
 $]$ $() =$
 $[$
 $(1, ()),$
 $]$

```

    <proof>

end

value [code] length (f4-punit DRLEX (map (λp. (p, ())) ((cyclic DRLEX 4)::(- ⇒0
rat) list)) ())

value [code] length (f4-punit DRLEX (map (λp. (p, ())) ((katsura DRLEX 2)::(-
⇒0 rat) list)) ())

end

```

18 Syzygies of Multivariate Polynomials

```

theory Syzygy
  imports Groebner-Bases More-MPoly-Type-Class
begin

```

In this theory we first introduce the general concept of *syzygies* in modules, and then provide a method for computing Gröbner bases of syzygy modules of lists of multivariate vector-polynomials. Since syzygies in this context are themselves represented by vector-polynomials, this method can be applied repeatedly to compute bases of syzygy modules of syzygies, and so on.

```

instance nat :: comm-powerprod <proof>

```

18.1 Syzygy Modules Generated by Sets

```

context module
begin

```

```

definition rep :: ('b ⇒0 'a) ⇒ 'b
  where rep r = (∑ v∈keys r. lookup r v * s v)

```

```

definition represents :: 'b set ⇒ ('b ⇒0 'a) ⇒ 'b ⇒ bool
  where represents B r x ⇔ (keys r ⊆ B ∧ local.rep r = x)

```

```

definition syzygy-module :: 'b set ⇒ ('b ⇒0 'a) set
  where syzygy-module B = {s. local.represents B s 0}

```

```

end

```

```

hide-const (open) real-vector.rep real-vector.represents real-vector.syzygy-module

```

```

context module
begin

```

```

lemma rep-monomial [simp]: rep (monomial c x) = c * s x
<proof>

```

lemma *rep-zero* [*simp*]: $\text{rep } 0 = 0$
<proof>

lemma *rep-uminus* [*simp*]: $\text{rep } (- r) = - \text{rep } r$
<proof>

lemma *rep-plus*: $\text{rep } (r + s) = \text{rep } r + \text{rep } s$
<proof>

lemma *rep-minus*: $\text{rep } (r - s) = \text{rep } r - \text{rep } s$
<proof>

lemma *rep-smult*: $\text{rep } (\text{monomial } c \ 0 \ * \ r) = c \ * \ \text{rep } r$
<proof>

lemma *rep-in-span*: $\text{rep } r \in \text{span } (\text{keys } r)$
<proof>

lemma *spanE-rep*:
 assumes $x \in \text{span } B$
 obtains r **where** $\text{keys } r \subseteq B$ **and** $x = \text{rep } r$
<proof>

lemma *representsI*:
 assumes $\text{keys } r \subseteq B$ **and** $\text{rep } r = x$
 shows $\text{represents } B \ r \ x$
<proof>

lemma *representsD1*:
 assumes $\text{represents } B \ r \ x$
 shows $\text{keys } r \subseteq B$
<proof>

lemma *representsD2*:
 assumes $\text{represents } B \ r \ x$
 shows $x = \text{rep } r$
<proof>

lemma *represents-mono*:
 assumes $\text{represents } B \ r \ x$ **and** $B \subseteq A$
 shows $\text{represents } A \ r \ x$
<proof>

lemma *represents-self*: $\text{represents } \{x\} \ (\text{monomial } 1 \ x) \ x$
<proof>

lemma *represents-zero*: $\text{represents } B \ 0 \ 0$
<proof>

lemma *represents-plus*:

assumes *represents* A r x **and** *represents* B s y

shows *represents* $(A \cup B)$ $(r + s)$ $(x + y)$

<proof>

lemma *represents-uminus*:

assumes *represents* B r x

shows *represents* B $(- r)$ $(- x)$

<proof>

lemma *represents-minus*:

assumes *represents* A r x **and** *represents* B s y

shows *represents* $(A \cup B)$ $(r - s)$ $(x - y)$

<proof>

lemma *represents-scale*:

assumes *represents* B r x

shows *represents* B (*monomial* c $0 * r$) $(c * s$ $x)$

<proof>

lemma *represents-in-span*:

assumes *represents* B r x

shows $x \in \text{span } B$

<proof>

lemma *syzygy-module-iff*: $s \in \text{syzygy-module } B \iff \text{represents } B$ s 0

<proof>

lemma *syzygy-moduleI*:

assumes *represents* B s 0

shows $s \in \text{syzygy-module } B$

<proof>

lemma *syzygy-moduleD*:

assumes $s \in \text{syzygy-module } B$

shows *represents* B s 0

<proof>

lemma *zero-in-syzygy-module*: $0 \in \text{syzygy-module } B$

<proof>

lemma *syzygy-module-closed-plus*:

assumes $s1 \in \text{syzygy-module } B$ **and** $s2 \in \text{syzygy-module } B$

shows $s1 + s2 \in \text{syzygy-module } B$

<proof>

lemma *syzygy-module-closed-minus*:

assumes $s1 \in \text{syzygy-module } B$ **and** $s2 \in \text{syzygy-module } B$

shows $s1 - s2 \in \text{syzygy-module } B$
 ⟨proof⟩

lemma *syzygy-module-closed-times-monomial*:
assumes $s \in \text{syzygy-module } B$
shows $\text{monomial } c \ 0 * s \in \text{syzygy-module } B$
 ⟨proof⟩

end

context *term-powerprod*
begin

lemma *keys-rep-subset*:
assumes $u \in \text{keys } (\text{pmdl.rep } r)$
obtains $t \ v$ **where** $t \in \text{Keys } (\text{Poly-Mapping.range } r)$ **and** $v \in \text{Keys } (\text{keys } r)$ **and**
 $u = t \oplus v$
 ⟨proof⟩

lemma *rep-mult-scalar*: $\text{pmdl.rep } (\text{punit.monom-mult } c \ 0 \ r) = c \odot \text{pmdl.rep } r$
 ⟨proof⟩

lemma *represents-mult-scalar*:
assumes $\text{pmdl.represents } B \ r \ x$
shows $\text{pmdl.represents } B \ (\text{punit.monom-mult } c \ 0 \ r) \ (c \odot x)$
 ⟨proof⟩

lemma *syzygy-module-closed-map-scale*: $s \in \text{pmdl.syzygy-module } B \implies c \cdot s \in \text{pmdl.syzygy-module } B$
 ⟨proof⟩

lemma *phull-syzygy-module*: $\text{phull } (\text{pmdl.syzygy-module } B) = \text{pmdl.syzygy-module } B$
 ⟨proof⟩

end

18.2 Polynomial Mappings on List-Indices

definition *pm-of-idx-pm* :: $('a \ \text{list}) \Rightarrow (\text{nat} \Rightarrow_0 'b) \Rightarrow 'a \Rightarrow_0 'b::\text{zero}$
where $\text{pm-of-idx-pm } xs \ f = \text{Abs-poly-mapping } (\lambda x. \text{lookup } f \ (\text{Min } \{i. i < \text{length } xs \wedge xs ! i = x\}))$ *when* $x \in \text{set } xs$

definition *idx-pm-of-pm* :: $('a \ \text{list}) \Rightarrow ('a \Rightarrow_0 'b) \Rightarrow \text{nat} \Rightarrow_0 'b::\text{zero}$
where $\text{idx-pm-of-pm } xs \ f = \text{Abs-poly-mapping } (\lambda i. \text{lookup } f \ (xs ! i))$ *when* $i < \text{length } xs$

lemma *lookup-pm-of-idx-pm*:
 $\text{lookup } (\text{pm-of-idx-pm } xs \ f) = (\lambda x. \text{lookup } f \ (\text{Min } \{i. i < \text{length } xs \wedge xs ! i = x\}))$

when $x \in \text{set } xs$)
(proof)

lemma *lookup-pm-of-idx-pm-distinct*:
assumes *distinct xs and $i < \text{length } xs$*
shows *lookup (pm-of-idx-pm xs f) (xs ! i) = lookup f i*
(proof)

lemma *keys-pm-of-idx-pm-subset*: *keys (pm-of-idx-pm xs f) \subseteq set xs*
(proof)

lemma *range-pm-of-idx-pm-subset*: *Poly-Mapping.range (pm-of-idx-pm xs f) \subseteq lookup f ‘ $\{0..<\text{length } xs\} - \{0\}$*
(proof)

corollary *range-pm-of-idx-pm-subset'*: *Poly-Mapping.range (pm-of-idx-pm xs f) \subseteq Poly-Mapping.range f*
(proof)

lemma *pm-of-idx-pm-zero [simp]*: *pm-of-idx-pm xs 0 = 0*
(proof)

lemma *pm-of-idx-pm-plus*: *pm-of-idx-pm xs (f + g) = pm-of-idx-pm xs f + pm-of-idx-pm xs g*
(proof)

lemma *pm-of-idx-pm-uminus*: *pm-of-idx-pm xs (- f) = - pm-of-idx-pm xs f*
(proof)

lemma *pm-of-idx-pm-minus*: *pm-of-idx-pm xs (f - g) = pm-of-idx-pm xs f - pm-of-idx-pm xs g*
(proof)

lemma *pm-of-idx-pm-monom-mult*: *pm-of-idx-pm xs (punit.monom-mult c 0 f) = punit.monom-mult c 0 (pm-of-idx-pm xs f)*
(proof)

lemma *pm-of-idx-pm-monomial*:
assumes *distinct xs*
shows *pm-of-idx-pm xs (monomial c i) = (monomial c (xs ! i) when $i < \text{length } xs$)*
(proof)

lemma *pm-of-idx-pm-take*:
assumes *keys f $\subseteq \{0..<j\}$*
shows *pm-of-idx-pm (take j xs) f = pm-of-idx-pm xs f*
(proof)

lemma *lookup-idx-pm-of-pm*: *lookup (idx-pm-of-pm xs f) = ($\lambda i.$ lookup f (xs ! i))*

when $i < \text{length } xs$
 $\langle \text{proof} \rangle$

lemma *keys-idx-pm-of-pm-subset*: $\text{keys } (\text{idx-pm-of-pm } xs \ f) \subseteq \{0..<\text{length } xs\}$
 $\langle \text{proof} \rangle$

lemma *idx-pm-of-pm-zero [simp]*: $\text{idx-pm-of-pm } xs \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma *idx-pm-of-pm-plus*: $\text{idx-pm-of-pm } xs \ (f + g) = \text{idx-pm-of-pm } xs \ f + \text{idx-pm-of-pm } xs \ g$
 $\langle \text{proof} \rangle$

lemma *idx-pm-of-pm-minus*: $\text{idx-pm-of-pm } xs \ (f - g) = \text{idx-pm-of-pm } xs \ f - \text{idx-pm-of-pm } xs \ g$
 $\langle \text{proof} \rangle$

lemma *pm-of-idx-pm-of-pm*:
assumes $\text{keys } f \subseteq \text{set } xs$
shows $\text{pm-of-idx-pm } xs \ (\text{idx-pm-of-pm } xs \ f) = f$
 $\langle \text{proof} \rangle$

lemma *idx-pm-of-pm-of-idx-pm*:
assumes *distinct* xs **and** $\text{keys } f \subseteq \{0..<\text{length } xs\}$
shows $\text{idx-pm-of-pm } xs \ (\text{pm-of-idx-pm } xs \ f) = f$
 $\langle \text{proof} \rangle$

18.3 POT Orders

context *ordered-term*
begin

definition *is-pot-ord* :: *bool*
where $\text{is-pot-ord} \longleftrightarrow (\forall u \ v. \text{component-of-term } u < \text{component-of-term } v \longrightarrow u \prec_t v)$

lemma *is-pot-ordI*:
assumes $\bigwedge u \ v. \text{component-of-term } u < \text{component-of-term } v \implies u \prec_t v$
shows *is-pot-ord*
 $\langle \text{proof} \rangle$

lemma *is-pot-ordD*:
assumes *is-pot-ord* **and** $\text{component-of-term } u < \text{component-of-term } v$
shows $u \prec_t v$
 $\langle \text{proof} \rangle$

lemma *is-pot-ordD2*:
assumes *is-pot-ord* **and** $u \preceq_t v$
shows $\text{component-of-term } u \leq \text{component-of-term } v$

<proof>

lemma *is-pot-ord*:

assumes *is-pot-ord*

shows $u \preceq_t v \iff (\text{component-of-term } u < \text{component-of-term } v \vee$
 $\text{component-of-term } u = \text{component-of-term } v \wedge \text{pp-of-term } u \preceq$
 $\text{pp-of-term } v)$ (**is** $?l \iff ?r$)

<proof>

definition *map-component* :: $('k \Rightarrow 'k) \Rightarrow 't \Rightarrow 't$

where *map-component* $f v = \text{term-of-pair } (\text{pp-of-term } v, f (\text{component-of-term } v))$

lemma *pair-of-map-component* [*term-simps*]:

pair-of-term (*map-component* $f v$) = (*pp-of-term* $v, f (\text{component-of-term } v)$)

<proof>

lemma *pp-of-map-component* [*term-simps*]: *pp-of-term* (*map-component* $f v$) = *pp-of-term* v

<proof>

lemma *component-of-map-component* [*term-simps*]:

component-of-term (*map-component* $f v$) = $f (\text{component-of-term } v)$

<proof>

lemma *map-component-term-of-pair* [*term-simps*]:

map-component $f (\text{term-of-pair } (t, k)) = \text{term-of-pair } (t, f k)$

<proof>

lemma *map-component-comp*: *map-component* $f (\text{map-component } g x) = \text{map-component } (\lambda k. f (g k)) x$

<proof>

lemma *map-component-id* [*term-simps*]: *map-component* $(\lambda k. k) x = x$

<proof>

lemma *map-component-inj*:

assumes *inj* f **and** *map-component* $f u = \text{map-component } f v$

shows $u = v$

<proof>

end

18.4 Gröbner Bases of Syzygy Modules

locale *gd-inf-term* =

gd-term *pair-of-term* *term-of-pair* *ord* *ord-strict* *ord-term* *ord-term-strict*

for *pair-of-term*:: $'t \Rightarrow ('a::\text{graded-dickson-powerprod} \times \text{nat})$

and *term-of-pair*:: $('a \times \text{nat}) \Rightarrow 't$

```

and ord::'a ⇒ 'a ⇒ bool (infixl ≤ 50)
and ord-strict (infixl < 50)
and ord-term::'t ⇒ 't ⇒ bool (infixl ≤t 50)
and ord-term-strict::'t ⇒ 't ⇒ bool (infixl <t 50)

```

begin

In order to compute a Gröbner basis of the syzygy module of a list bs of polynomials, one first needs to “lift” bs to a new list bs' by adding further components, compute a Gröbner basis gs of bs' , and then filter out those elements of gs whose only non-zero components are those that were newly added to bs . Function *init-syzygy-list* takes care of constructing bs' , and function *filter-syzygy-basis* does the filtering. Function *proj-orig-basis*, finally, projects the Gröbner basis gs of bs' to a Gröbner basis of the original list bs .

```

definition lift-poly-syz :: nat ⇒ (t ⇒0 'b) ⇒ nat ⇒ (t ⇒0 'b::semiring-1)
  where lift-poly-syz n b i = Abs-poly-mapping
    (λx. if pair-of-term x = (0, i) then 1
      else if n ≤ component-of-term x then lookup b (map-component (λk.
k - n) x)
      else 0)

```

```

definition proj-poly-syz :: nat ⇒ (t ⇒0 'b) ⇒ (t ⇒0 'b::semiring-1)
  where proj-poly-syz n b = Poly-Mapping.map-key (λx. map-component (λk. k +
n) x) b

```

```

definition cofactor-list-syz :: nat ⇒ (t ⇒0 'b) ⇒ ('a ⇒0 'b::semiring-1) list
  where cofactor-list-syz n b = map (λi. proj-poly i b) [0..n]

```

```

definition init-syzygy-list :: (t ⇒0 'b) list ⇒ (t ⇒0 'b::semiring-1) list
  where init-syzygy-list bs = map-idx (lift-poly-syz (length bs)) bs 0

```

```

definition proj-orig-basis :: nat ⇒ (t ⇒0 'b) list ⇒ (t ⇒0 'b::semiring-1) list
  where proj-orig-basis n bs = map (proj-poly-syz n) bs

```

```

definition filter-syzygy-basis :: nat ⇒ (t ⇒0 'b) list ⇒ (t ⇒0 'b::semiring-1) list
  where filter-syzygy-basis n bs = [b←bs. component-of-term ‘ keys b ⊆ {0..n}]

```

```

definition syzygy-module-list :: (t ⇒0 'b) list ⇒ (t ⇒0 'b::comm-ring-1) set
  where syzygy-module-list bs = atomize-poly ‘ idx-pm-of-pm bs ‘ pmdl.syzygy-module
(set bs)

```

18.4.1 lift-poly-syz

lemma *keys-lift-poly-syz-aux*:

```

{x. (if pair-of-term x = (0, i) then 1
  else if n ≤ component-of-term x then lookup b (map-component (λk. k - n)
x)
  else 0) ≠ 0} ⊆ insert (term-of-pair (0, i)) (map-component (λk. k + n) ‘
keys b)

```

(is ?l \subseteq ?r) for b::t \Rightarrow_0 'b::semiring-1
 ⟨proof⟩

lemma lookup-lift-poly-syz:

lookup (lift-poly-syz n b i) =
 (λx . if pair-of-term x = (0, i) then 1 else if $n \leq$ component-of-term x then
 lookup b (map-component (λk . k - n) x) else 0)
 ⟨proof⟩

corollary lookup-lift-poly-syz-alt:

lookup (lift-poly-syz n b i) (term-of-pair (t, j)) =
 (if (t, j) = (0, i) then 1 else if $n \leq j$ then lookup b (term-of-pair (t, j -
 n)) else 0)
 ⟨proof⟩

lemma keys-lift-poly-syz:

keys (lift-poly-syz n b i) = insert (term-of-pair (0, i)) (map-component (λk . k +
 n) ' keys b)
 ⟨proof⟩

18.4.2 proj-poly-syz

lemma inj-map-component-plus: inj (map-component (λk . k + n))
 ⟨proof⟩

lemma lookup-proj-poly-syz: lookup (proj-poly-syz n p) x = lookup p (map-component
 (λk . k + n) x)
 ⟨proof⟩

lemma lookup-proj-poly-syz-alt:

lookup (proj-poly-syz n p) (term-of-pair (t, i)) = lookup p (term-of-pair (t, i +
 n))
 ⟨proof⟩

lemma keys-proj-poly-syz: keys (proj-poly-syz n p) = map-component (λk . k + n)
 - ' keys p
 ⟨proof⟩

lemma proj-poly-syz-zero [simp]: proj-poly-syz n 0 = 0
 ⟨proof⟩

lemma proj-poly-syz-plus: proj-poly-syz n (p + q) = proj-poly-syz n p + proj-poly-syz
 n q
 ⟨proof⟩

lemma proj-poly-syz-sum: proj-poly-syz n (sum f A) = ($\sum a \in A$. proj-poly-syz n (f
 a))
 ⟨proof⟩

lemma *proj-poly-syz-sum-list*: $\text{proj-poly-syz } n \text{ (sum-list } xs) = \text{sum-list (map (proj-poly-syz } n) xs)$

<proof>

lemma *proj-poly-syz-monom-mult*:

$\text{proj-poly-syz } n \text{ (monom-mult } c \ t \ p) = \text{monom-mult } c \ t \ (\text{proj-poly-syz } n \ p)$

<proof>

lemma *proj-poly-syz-mult-scalar*:

$\text{proj-poly-syz } n \text{ (mult-scalar } q \ p) = \text{mult-scalar } q \ (\text{proj-poly-syz } n \ p)$

<proof>

lemma *proj-poly-syz-lift-poly-syz*:

assumes $i < n$

shows $\text{proj-poly-syz } n \text{ (lift-poly-syz } n \ p \ i) = p$

<proof>

lemma *proj-poly-syz-eq-zero-iff*: $\text{proj-poly-syz } n \ p = 0 \longleftrightarrow (\text{component-of-term } \text{ `keys } p \subseteq \{0..<n\})$

<proof>

lemma *component-of-lt-ge*:

assumes *is-pot-ord* **and** $\text{proj-poly-syz } n \ p \neq 0$

shows $n \leq \text{component-of-term (lt } p)$

<proof>

lemma *lt-proj-poly-syz*:

assumes *is-pot-ord* **and** $\text{proj-poly-syz } n \ p \neq 0$

shows $\text{lt (proj-poly-syz } n \ p) = \text{map-component } (\lambda k. k - n) \text{ (lt } p) \text{ (is - = ?l)}$

<proof>

lemma *proj-proj-poly-syz*: $\text{proj-poly } k \ (\text{proj-poly-syz } n \ p) = \text{proj-poly } (k + n) \ p$

<proof>

lemma *poly-mapping-eqI-proj-syz*:

assumes $\text{proj-poly-syz } n \ p = \text{proj-poly-syz } n \ q$

and $\bigwedge k. k < n \implies \text{proj-poly } k \ p = \text{proj-poly } k \ q$

shows $p = q$

<proof>

18.4.3 cofactor-list-syz

lemma *length-cofactor-list-syz [simp]*: $\text{length (cofactor-list-syz } n \ p) = n$

<proof>

lemma *cofactor-list-syz-nth*:

assumes $i < n$

shows $(\text{cofactor-list-syz } n \ p) ! i = \text{proj-poly } i \ p$

<proof>

lemma *cofactor-list-syz-zero* [simp]: $\text{cofactor-list-syz } n \ 0 = \text{replicate } n \ 0$
 ⟨proof⟩

lemma *cofactor-list-syz-plus*:
 $\text{cofactor-list-syz } n \ (p + q) = \text{map2 } (+) \ (\text{cofactor-list-syz } n \ p) \ (\text{cofactor-list-syz } n \ q)$
 ⟨proof⟩

18.4.4 *init-syzygy-list*

lemma *length-init-syzygy-list* [simp]: $\text{length } (\text{init-syzygy-list } bs) = \text{length } bs$
 ⟨proof⟩

lemma *init-syzygy-list-nth*:
assumes $i < \text{length } bs$
shows $(\text{init-syzygy-list } bs) ! i = \text{lift-poly-syz } (\text{length } bs) \ (bs ! i)$
 ⟨proof⟩

lemma *Keys-init-syzygy-list*:
 $\text{Keys } (\text{set } (\text{init-syzygy-list } bs)) =$
 $\text{map-component } (\lambda k. k + \text{length } bs) \ ' \ \text{Keys } (\text{set } bs) \cup \ (\lambda i. \text{term-of-pair } (0, i))$
 $\ ' \ \{0..<\text{length } bs\}$
 ⟨proof⟩

lemma *pp-of-Keys-init-syzygy-list-subset*:
 $\text{pp-of-term } ' \ \text{Keys } (\text{set } (\text{init-syzygy-list } bs)) \subseteq \text{insert } 0 \ (\text{pp-of-term } ' \ \text{Keys } (\text{set } bs))$
 ⟨proof⟩

lemma *pp-of-Keys-init-syzygy-list-superset*:
 $\text{pp-of-term } ' \ \text{Keys } (\text{set } bs) \subseteq \text{pp-of-term } ' \ \text{Keys } (\text{set } (\text{init-syzygy-list } bs))$
 ⟨proof⟩

lemma *pp-of-Keys-init-syzygy-list*:
assumes $bs \neq []$
shows $\text{pp-of-term } ' \ \text{Keys } (\text{set } (\text{init-syzygy-list } bs)) = \text{insert } 0 \ (\text{pp-of-term } ' \ \text{Keys } (\text{set } bs))$
 ⟨proof⟩

lemma *component-of-Keys-init-syzygy-list*:
 $\text{component-of-term } ' \ \text{Keys } (\text{set } (\text{init-syzygy-list } bs)) =$
 $(+) \ (\text{length } bs) \ ' \ \text{component-of-term } ' \ \text{Keys } (\text{set } bs) \cup \ \{0..<\text{length } bs\}$
 ⟨proof⟩

lemma *proj-lift-poly-syz*:
assumes $j < n$
shows $\text{proj-poly } j \ (\text{lift-poly-syz } n \ p \ i) = (1 \ \text{when } j = i)$
 ⟨proof⟩

18.4.5 *proj-orig-basis*

lemma *length-proj-orig-basis* [simp]: $\text{length } (\text{proj-orig-basis } n \text{ } bs) = \text{length } bs$
<proof>

lemma *proj-orig-basis-nth*:
 assumes $i < \text{length } bs$
 shows $(\text{proj-orig-basis } n \text{ } bs) ! i = \text{proj-poly-syz } n \text{ } (bs ! i)$
<proof>

lemma *proj-orig-basis-init-syzygy-list* [simp]:
 $\text{proj-orig-basis } (\text{length } bs) \text{ } (\text{init-syzygy-list } bs) = bs$
<proof>

lemma *set-proj-orig-basis*: $\text{set } (\text{proj-orig-basis } n \text{ } bs) = \text{proj-poly-syz } n \text{ } \text{' } \text{set } bs$
<proof>

The following lemma could be generalized from *proj-poly-syz* to arbitrary module homomorphisms, i. e. functions respecting 0, addition and scalar multiplication.

lemma *pmdl-proj-orig-basis'*:
 $\text{pmdl } (\text{set } (\text{proj-orig-basis } n \text{ } bs)) = \text{proj-poly-syz } n \text{ } \text{' } \text{pmdl } (\text{set } bs) \text{ } (\text{is } ?A = ?B)$
<proof>

18.4.6 *filter-syzygy-basis*

lemma *filter-syzygy-basis-alt*: $\text{filter-syzygy-basis } n \text{ } bs = [b \leftarrow bs. \text{proj-poly-syz } n \text{ } b = 0]$
<proof>

lemma *set-filter-syzygy-basis*:
 $\text{set } (\text{filter-syzygy-basis } n \text{ } bs) = \{b \in \text{set } bs. \text{proj-poly-syz } n \text{ } b = 0\}$
<proof>

18.4.7 *syzygy-module-list*

lemma *syzygy-module-listI*:
 assumes $s' \in \text{pmdl.syzygy-module } (\text{set } bs)$ **and** $s = \text{atomize-poly } (\text{idx-pm-of-pm } bs \text{ } s')$
 shows $s \in \text{syzygy-module-list } bs$
<proof>

lemma *syzygy-module-listE*:
 assumes $s \in \text{syzygy-module-list } bs$
 obtains s' **where** $s' \in \text{pmdl.syzygy-module } (\text{set } bs)$ **and** $s = \text{atomize-poly } (\text{idx-pm-of-pm } bs \text{ } s')$
<proof>

lemma *monom-mult-atomize*:

$monom-mult\ c\ t\ (atomize-poly\ p) = atomize-poly\ (MPoly-Type-Class.punit.monom-mult\ (monomial\ c\ t)\ 0\ p)$
 ⟨proof⟩

lemma *punit-monom-mult-monomial-idx-pm-of-pm*:

$MPoly-Type-Class.punit.monom-mult\ (monomial\ c\ t)\ (0::nat)\ (idx-pm-of-pm\ bs\ s) =$
 $idx-pm-of-pm\ bs\ (MPoly-Type-Class.punit.monom-mult\ (monomial\ c\ t)\ (0::'t\ \Rightarrow_0\ 'b::ring-1)\ s)$
 ⟨proof⟩

lemma *syzygy-module-list-closed-monom-mult*:

assumes $s \in syzygy-module-list\ bs$
shows $monom-mult\ c\ t\ s \in syzygy-module-list\ bs$
 ⟨proof⟩

lemma *pm-dl-syzygy-module-list [simp]*: $pm-dl\ (syzygy-module-list\ bs) = syzygy-module-list\ bs$
 ⟨proof⟩

The following lemma also holds without the distinctness constraint on bs , but then the proof becomes more difficult.

lemma *syzygy-module-listI'*:

assumes *distinct* bs **and** *sum-list* $(map2\ mult-scalar\ (cofactor-list-syz\ (length\ bs)\ s)\ bs) = 0$
and *component-of-term* $\text{'keys } s \subseteq \{0..<length\ bs\}$
shows $s \in syzygy-module-list\ bs$
 ⟨proof⟩

lemma *component-of-syzygy-module-list*:

assumes $s \in syzygy-module-list\ bs$
shows *component-of-term* $\text{'keys } s \subseteq \{0..<length\ bs\}$
 ⟨proof⟩

lemma *map2-mult-scalar-proj-poly-syz*:

$map2\ mult-scalar\ xs\ (map\ (proj-poly-syz\ n)\ ys) =$
 $map\ (proj-poly-syz\ n\ \circ\ (\lambda(x, y).\ mult-scalar\ x\ y))\ (zip\ xs\ ys)$
 ⟨proof⟩

lemma *map2-times-proj*:

$map2\ (*)\ xs\ (map\ (proj-poly\ k)\ ys) = map\ (proj-poly\ k\ \circ\ (\lambda(x, y).\ x\ \odot\ y))\ (zip\ xs\ ys)$
 ⟨proof⟩

Probably the following lemma also holds without the distinctness constraint on bs .

lemma *syzygy-module-list-subset*:

assumes *distinct* bs
shows $syzygy-module-list\ bs \subseteq pm-dl\ (set\ (init-syzygy-list\ bs))$

<proof>

18.4.8 Cofactors

lemma *map2-mult-scalar-plus*:

$map2 (\odot) (map2 (+) xs ys) zs = map2 (+) (map2 (\odot) xs zs) (map2 (\odot) ys zs)$
<proof>

lemma *syz-cofactors*:

assumes $p \in pmdl (set (init-syzygy-list bs))$
shows $proj-poly-syz (length bs) p = sum-list (map2 mult-scalar (cofactor-list-syz (length bs) p) bs)$
<proof>

18.4.9 Modules

lemma *pmdl-proj-orig-basis*:

assumes $pmdl (set gs) = pmdl (set (init-syzygy-list bs))$
shows $pmdl (set (proj-orig-basis (length bs) gs)) = pmdl (set bs)$
<proof>

lemma *pmdl-filter-syzygy-basis-subset*:

assumes *distinct bs* **and** $pmdl (set gs) = pmdl (set (init-syzygy-list bs))$
shows $pmdl (set (filter-syzygy-basis (length bs) gs)) \subseteq pmdl (syzygy-module-list bs)$
<proof>

lemma *ex-filter-syzygy-basis-adds-lt*:

assumes *is-pot-ord* **and** *distinct bs* **and** *is-Groebner-basis (set gs)*
and $pmdl (set gs) = pmdl (set (init-syzygy-list bs))$
and $f \in pmdl (syzygy-module-list bs)$ **and** $f \neq 0$
shows $\exists g \in set (filter-syzygy-basis (length bs) gs). g \neq 0 \wedge lt g adds_t lt f$
<proof>

lemma *pmdl-filter-syzygy-basis*:

fixes $bs::('t \Rightarrow_0 'b::field) list$
assumes *is-pot-ord* **and** *distinct bs* **and** *is-Groebner-basis (set gs)* **and**
 $pmdl (set gs) = pmdl (set (init-syzygy-list bs))$
shows $pmdl (set (filter-syzygy-basis (length bs) gs)) = syzygy-module-list bs$
<proof>

18.4.10 Gröbner Bases

lemma *proj-orig-basis-isGB*:

assumes *is-pot-ord* **and** *is-Groebner-basis (set gs)* **and** $pmdl (set gs) = pmdl (set (init-syzygy-list bs))$
shows *is-Groebner-basis (set (proj-orig-basis (length bs) gs))*
<proof>

lemma *filter-syzygy-basis-isGB*:

```

assumes is-pot-ord and distinct bs and is-Groebner-basis (set gs)
and pmdl (set gs) = pmdl (set (init-syzygy-list bs))
shows is-Groebner-basis (set (filter-syzygy-basis (length bs) gs))
  <proof>

```

end

end

19 Sample Computations of Syzygies

theory *Syzygy-Examples*

imports *Buchberger Algorithm-Schema-Impl Syzygy Code-Target-Rat*
begin

19.1 Preparations

We must define the following four constants outside the global interpretation, since otherwise their types are too general.

definition *splus-pprod* :: ('a::nat, 'b::nat) pp \Rightarrow -
where *splus-pprod* = *pprod.splus*

definition *monom-mult-pprod* :: 'c::semiring-0 \Rightarrow ('a::nat, 'b::nat) pp \Rightarrow (((('a, 'b) pp \times nat) \Rightarrow_0 'c) \Rightarrow -
where *monom-mult-pprod* = *pprod.monom-mult*

definition *mult-scalar-pprod* :: (('a::nat, 'b::nat) pp \Rightarrow_0 'c::semiring-0) \Rightarrow (((('a, 'b) pp \times nat) \Rightarrow_0 'c) \Rightarrow -
where *mult-scalar-pprod* = *pprod.mult-scalar*

definition *adds-term-pprod* :: (('a::nat, 'b::nat) pp \times -) \Rightarrow -
where *adds-term-pprod* = *pprod.adds-term*

lemma (**in** *gd-term*) *compute-trd-aux* [*code*]:

```

\Rightarrow trd-aux fs (tail p) (plus-monomial-less r (lc p) (lt p))
    | Some f  $\Rightarrow$  trd-aux fs (tail p - monom-mult (lc p / lc f) (lp p - lp f) (tail
f)) r
  )
  <proof>

```

locale *gd-nat-inf-term* = *gd-nat-term pair-of-term term-of-pair cmp-term*

for *pair-of-term*::'t::nat-term \Rightarrow ('a::{nat-term,graded-dickson-powerprod} \times nat)

and *term-of-pair*::('a × nat) ⇒ 't
and *cmp-term*
begin

sublocale *aux*: *gd-inf-term pair-of-term term-of-pair*
 $\lambda s t. \text{le-of-nat-term-order } \text{cmp-term } (\text{term-of-pair } (s, \text{the-min})) (\text{term-of-pair } (t, \text{the-min}))$
 $\lambda s t. \text{lt-of-nat-term-order } \text{cmp-term } (\text{term-of-pair } (s, \text{the-min})) (\text{term-of-pair } (t, \text{the-min}))$
le-of-nat-term-order cmp-term
lt-of-nat-term-order cmp-term ⟨*proof*⟩

definition *lift-keys* :: nat ⇒ ('t, 'b) *oalist-ntm* ⇒ ('t, 'b::*semiring-0*) *oalist-ntm*
where *lift-keys* *i xs* = *oalist-of-list-ntm* (*map-raw* ($\lambda kv. (\text{map-component } ((+) i) (\text{fst } kv), \text{snd } kv)$) (*list-of-oalist-ntm* *xs*))

lemma *list-of-oalist-lift-keys*:
 $\text{list-of-oalist-ntm } (\text{lift-keys } i \text{ } xs) = (\text{map-raw } (\lambda kv. (\text{map-component } ((+) i) (\text{fst } kv), \text{snd } kv)) (\text{list-of-oalist-ntm } xs))$
 ⟨*proof*⟩

Regardless of whether the above lemma holds (which might be the case) or not, we can use *lift-keys* in computations. Now, however, it is implemented rather inefficiently, because the list resulting from the application of *map-raw* is sorted again. That should not be a big problem though, since *lift-keys* is applied only once to every input polynomial before computing syzygies.

lemma *lookup-lift-keys-plus*:
 $\text{lookup } (\text{MP-oalist } (\text{lift-keys } i \text{ } xs)) (\text{term-of-pair } (t, i + k)) = \text{lookup } (\text{MP-oalist } xs) (\text{term-of-pair } (t, k))$
 (is ?l = ?r)
 ⟨*proof*⟩

lemma *keys-lift-keys-subset*:
 $\text{keys } (\text{MP-oalist } (\text{lift-keys } i \text{ } xs)) \subseteq (\text{map-component } ((+) i) \text{ } \text{'keys } (\text{MP-oalist } xs))$
 (is ?l ⊆ ?r)
 ⟨*proof*⟩

end

global-interpretation *pprod'*: *gd-nat-inf-term* $\lambda x::('a, 'b) \text{ } pp \times \text{nat. } x \lambda x. x \text{ } \text{cmp-term}$
rewrites *pprod.pp-of-term* = *fst*
and *pprod.component-of-term* = *snd*
and *pprod.splus* = *splus-pprod*
and *pprod.monom-mult* = *monom-mult-pprod*
and *pprod.mult-scalar* = *mult-scalar-pprod*
and *pprod.adds-term* = *adds-term-pprod*
for *cmp-term* :: (('a::nat, 'b::nat) *pp* × nat) *nat-term-order*
defines *shift-map-keys-pprod* = *pprod'.shift-map-keys*
and *lift-keys-pprod* = *pprod'.lift-keys*

```

and min-term-pprod = pprod'.min-term
and lt-pprod = pprod'.lt
and lc-pprod = pprod'.lc
and tail-pprod = pprod'.tail
and comp-opt-p-pprod = pprod'.comp-opt-p
and ord-p-pprod = pprod'.ord-p
and ord-strict-p-pprod = pprod'.ord-strict-p
and find-adds-pprod = pprod'.find-adds
and trd-aux-pprod = pprod'.trd-aux
and trd-pprod = pprod'.trd
and spoly-pprod = pprod'.spoly
and count-const-lt-components-pprod = pprod'.count-const-lt-components
and count-rem-components-pprod = pprod'.count-rem-components
and const-lt-component-pprod = pprod'.const-lt-component
and full-gb-pprod = pprod'.full-gb
and keys-to-list-pprod = pprod'.keys-to-list
and Keys-to-list-pprod = pprod'.Keys-to-list
and add-pairs-single-sorted-pprod = pprod'.add-pairs-single-sorted
and add-pairs-pprod = pprod'.add-pairs
and canon-pair-order-aux-pprod = pprod'.canon-pair-order-aux
and canon-basis-order-pprod = pprod'.canon-basis-order
and new-pairs-sorted-pprod = pprod'.new-pairs-sorted
and component-crit-pprod = pprod'.component-crit
and chain-ncrit-pprod = pprod'.chain-ncrit
and chain-ocrit-pprod = pprod'.chain-ocrit
and apply-icrit-pprod = pprod'.apply-icrit
and apply-ncrit-pprod = pprod'.apply-ncrit
and apply-ocrit-pprod = pprod'.apply-ocrit
and trdsp-pprod = pprod'.trdsp
and gb-sel-pprod = pprod'.gb-sel
and gb-red-aux-pprod = pprod'.gb-red-aux
and gb-red-pprod = pprod'.gb-red
and gb-aux-pprod = pprod'.gb-aux
and gb-pprod = pprod'.gb
and filter-syzygy-basis-pprod = pprod'.aux.filter-syzygy-basis
and init-syzygy-list-pprod = pprod'.aux.init-syzygy-list
and lift-poly-syz-pprod = pprod'.aux.lift-poly-syz
and map-component-pprod = pprod'.map-component
⟨proof⟩

```

lemma *compute-adds-term-pprod* [code]:
 $adds-term-pprod\ u\ v = (snd\ u = snd\ v \wedge adds-pp-add-linorder\ (fst\ u)\ (fst\ v))$
⟨proof⟩

lemma *compute-splus-pprod* [code]: $splus-pprod\ t\ (s,\ i) = (t + s,\ i)$
⟨proof⟩

lemma *compute-shift-map-keys-pprod* [code abstract]:
 $list-of-oalist-ntm\ (shift-map-keys-pprod\ t\ f\ xs) = map-raw\ (\lambda(k,\ v).\ (splus-pprod$

$t\ k, f\ v)$ (*list-of-oalist-ntm* xs)
 ⟨*proof*⟩

lemma *compute-trd-pprod* [*code*]: *trd-pprod to fs p = trd-aux-pprod to fs p (change-ord to 0)*
 ⟨*proof*⟩

lemmas [*code*] = *conversep-iff*

lemma *POT-is-pot-ord: pprod'.is-pot-ord (TYPE('a::nat)) (TYPE('b::nat)) (POT to)*
 ⟨*proof*⟩

definition $Vec_0 :: nat \Rightarrow (('a, nat) pp \Rightarrow_0 'b) \Rightarrow (('a::nat, nat) pp \times nat) \Rightarrow_0 'b::semiring-1$ **where**
 $Vec_0\ i\ p = mult\ scalar\ pprod\ p\ (Poly\ Mapping.single\ (0, i)\ 1)$

definition *syzygy-basis to bs =*
 $filter\ syzygy\ basis\ pprod\ (length\ bs)\ (map\ fst\ (gb\ pprod\ (POT\ to)\ (map\ (\lambda p. (p, ()))\ (init\ syzygy\ list\ pprod\ bs))\ ()))$

thm *pprod'.aux.filter-syzygy-basis-isGB[OF POT-is-pot-ord]*

lemma *lift-poly-syz-MP-oalist* [*code*]:
 $lift\ poly\ syz\ pprod\ n\ (MP\ oalist\ xs)\ i = MP\ oalist\ (OAlist\ insert\ ntm\ ((0, i), 1)\ (lift\ keys\ pprod\ n\ xs))$
 ⟨*proof*⟩

19.2 Computations

experiment begin interpretation *trivariate₀-rat* ⟨*proof*⟩

lemma
 $syzygy\ basis\ DRLEX\ [Vec_0\ 0\ (X^2 * Z^3 + 3 * X^2 * Y), Vec_0\ 0\ (X * Y * Z + 2 * Y^2)] =$
 $[Vec_0\ 0\ (C_0\ (1 / 3) * X * Y * Z + C_0\ (2 / 3) * Y^2) + Vec_0\ 1\ (C_0\ (-1 / 3) * X^2 * Z^3 - X^2 * Y)]$
 ⟨*proof*⟩

value [*code*] *syzygy-basis DRLEX* [$Vec_0\ 0\ (X^2 * Z^3 + 3 * X^2 * Y)$, $Vec_0\ 0\ (X * Y * Z + 2 * Y^2)$, $Vec_0\ 0\ (X - Y + 3 * Z)$]

lemma
 $map\ fst\ (gb\ pprod\ (POT\ DRLEX)\ (map\ (\lambda p. (p, ()))\ (init\ syzygy\ list\ pprod\ [Vec_0\ 0\ (X^4 + 3 * X^2 * Y), Vec_0\ 0\ (Y^3 + 2 * X * Z), Vec_0\ 0\ (Z^2 - X - Y)]))\ ())) =$
 [
 $Vec_0\ 0\ 1 + Vec_0\ 3\ (X^4 + 3 * X^2 * Y),$
 $Vec_0\ 1\ 1 + Vec_0\ 3\ (Y^3 + 2 * X * Z),$

$$\begin{aligned} & \text{Vec}_0 0 (Y^{\wedge} 3 + 2 * X * Z) - \text{Vec}_0 1 (X^{\wedge} 4 + 3 * X^2 * Y), \\ & \text{Vec}_0 2 1 + \text{Vec}_0 3 (Z^2 - X - Y), \\ & \text{Vec}_0 1 (Z^2 - X - Y) - \text{Vec}_0 2 (Y^{\wedge} 3 + 2 * X * Z), \\ & \text{Vec}_0 0 (Z^2 - X - Y) - \text{Vec}_0 2 (X^{\wedge} 4 + 3 * X^2 * Y), \\ & \text{Vec}_0 0 (- (Y^{\wedge} 3 * Z^2) + Y^{\wedge} 4 + X * Y^{\wedge} 3 + 2 * X^2 * Z + 2 * X * Y * \\ & Z - 2 * X * Z^{\wedge} 3) + \\ & \text{Vec}_0 1 (X^{\wedge} 4 * Z^2 - X^{\wedge} 5 - X^{\wedge} 4 * Y - 3 * X^{\wedge} 3 * Y - 3 * X^2 * Y^2 \\ & + 3 * X^2 * Y * Z^2) \\ &] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma

$$\begin{aligned} & \text{syzygy-basis DRLEX} [\text{Vec}_0 0 (X^{\wedge} 4 + 3 * X^2 * Y), \text{Vec}_0 0 (Y^{\wedge} 3 + 2 * X * \\ & Z), \text{Vec}_0 0 (Z^2 - X - Y)] = \\ & [\\ & \text{Vec}_0 0 (Y^{\wedge} 3 + 2 * X * Z) - \text{Vec}_0 1 (X^{\wedge} 4 + 3 * X^2 * Y), \\ & \text{Vec}_0 1 (Z^2 - X - Y) - \text{Vec}_0 2 (Y^{\wedge} 3 + 2 * X * Z), \\ & \text{Vec}_0 0 (Z^2 - X - Y) - \text{Vec}_0 2 (X^{\wedge} 4 + 3 * X^2 * Y), \\ & \text{Vec}_0 0 (- (Y^{\wedge} 3 * Z^2) + Y^{\wedge} 4 + X * Y^{\wedge} 3 + 2 * X^2 * Z + 2 * X * Y * \\ & Z - 2 * X * Z^{\wedge} 3) + \\ & \text{Vec}_0 1 (X^{\wedge} 4 * Z^2 - X^{\wedge} 5 - X^{\wedge} 4 * Y - 3 * X^{\wedge} 3 * Y - 3 * X^2 * Y^2 \\ & + 3 * X^2 * Y * Z^2) \\ &] \\ & \langle \text{proof} \rangle \end{aligned}$$

value [code] $\text{syzygy-basis DRLEX} [\text{Vec}_0 0 (X * Y - Z), \text{Vec}_0 0 (X * Z - Y), \text{Vec}_0 0 (Y * Z - X)]$

lemma

$$\begin{aligned} & \text{map fst (gb-pprod (POT DRLEX) (map (\lambda p. (p, ())) (init-syzygy-list-pprod \\ & [\text{Vec}_0 0 (X * Y - Z), \text{Vec}_0 0 (X * Z - Y), \text{Vec}_0 0 (Y * Z - X)])) ())) = \\ & [\\ & \text{Vec}_0 0 1 + \text{Vec}_0 3 (X * Y - Z), \\ & \text{Vec}_0 1 1 + \text{Vec}_0 3 (X * Z - Y), \\ & \text{Vec}_0 2 1 + \text{Vec}_0 3 (Y * Z - X), \\ & \text{Vec}_0 0 (- X * Z + Y) + \text{Vec}_0 1 (X * Y - Z), \\ & \text{Vec}_0 0 (- Y * Z + X) + \text{Vec}_0 2 (X * Y - Z), \\ & \text{Vec}_0 1 (- Y * Z + X) + \text{Vec}_0 2 (X * Z - Y), \\ & \text{Vec}_0 1 (- Y) + \text{Vec}_0 2 (X) + \text{Vec}_0 3 (Y^{\wedge} 2 - X^{\wedge} 2), \\ & \text{Vec}_0 0 (Z) + \text{Vec}_0 2 (-X) + \text{Vec}_0 3 (X^{\wedge} 2 - Z^{\wedge} 2), \\ & \text{Vec}_0 0 (Y - Y * Z^{\wedge} 2) + \text{Vec}_0 1 (Y^{\wedge} 2 * Z - Z) + \text{Vec}_0 2 (Y^{\wedge} 2 - Z^{\wedge} \\ & 2), \\ & \text{Vec}_0 0 (- Y) + \text{Vec}_0 1 (- (X * Y)) + \text{Vec}_0 2 (X^{\wedge} 2 - 1) + \text{Vec}_0 3 (X - \\ & X^{\wedge} 3) \\ &] \\ & \langle \text{proof} \rangle \end{aligned}$$

lemma

$$\text{syzygy-basis DRLEX} [\text{Vec}_0 0 (X * Y - Z), \text{Vec}_0 0 (X * Z - Y), \text{Vec}_0 0 (Y * Z - X)]$$


```

Z - X)] =
[
  Vec0 0 (- X * Z + Y) + Vec0 1 (X * Y - Z),
  Vec0 0 (- Y * Z + X) + Vec0 2 (X * Y - Z),
  Vec0 1 (- Y * Z + X) + Vec0 2 (X * Z - Y),
  Vec0 0 (Y - Y * Z ^ 2) + Vec0 1 (Y ^ 2 * Z - Z) + Vec0 2 (Y ^ 2 - Z ^
2)
]
⟨proof⟩

```

end

end

theory Groebner-PM

imports Polynomials.MPoly-PM Reduced-GB

begin

We prove results that hold specifically for Gröbner bases in polynomial rings, where the polynomials really have *indeterminates*.

context pm-powerprod

begin

lemmas finite-reduced-GB-Polys =

punit.finite-reduced-GB-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-is-reduced-GB-Polys =

punit.reduced-GB-is-reduced-GB-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-is-GB-Polys =

punit.reduced-GB-is-GB-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-is-auto-reduced-Polys =

punit.reduced-GB-is-auto-reduced-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-is-monic-set-Polys =

punit.reduced-GB-is-monic-set-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-nonzero-Polys =

punit.reduced-GB-nonzero-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-ideal-Polys =

punit.reduced-GB-pmdl-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-unique-Polys =

punit.reduced-GB-unique-dgrad-p-set[simplified, OF dickson-grading-varnum, where m=0, simplified dgrad-p-set-varnum]

lemmas reduced-GB-Polys =

punit.reduced-GB-dgrad-p-set[simplified, OF dickson-grading-varnum, where $m=0$, simplified dgrad-p-set-varnum]
lemmas *ideal-eq-UNIV-iff-reduced-GB-eq-one-Polys =*
ideal-eq-UNIV-iff-reduced-GB-eq-one-dgrad-p-set[simplified, OF dickson-grading-varnum,
where $m=0$, simplified dgrad-p-set-varnum]

19.3 Univariate Polynomials

lemma (*in -*) *adds-univariate-linear:*
assumes *finite X and card X ≤ 1 and $s \in .[X]$ and $t \in .[X]$*
obtains *s adds t | t adds s*
<proof>

context
fixes *X :: 'x set*
assumes *fin-X: finite X and card-X: card X ≤ 1*
begin

lemma *ord-iff-adds-univariate:*
assumes *$s \in .[X]$ and $t \in .[X]$*
shows *$s \preceq t \iff s$ adds t*
<proof>

lemma *adds-iff-deg-le-univariate:*
assumes *$s \in .[X]$ and $t \in .[X]$*
shows *s adds $t \iff \text{deg-pm } s \leq \text{deg-pm } t$*
<proof>

corollary *ord-iff-deg-le-univariate: $s \in .[X] \implies t \in .[X] \implies s \preceq t \iff \text{deg-pm } s$*
 $\leq \text{deg-pm } t$
<proof>

lemma *poly-deg-univariate:*
assumes *$p \in P[X]$*
shows *$\text{poly-deg } p = \text{deg-pm } (\text{lpp } p)$*
<proof>

lemma *reduced-GB-univariate-cases:*
assumes *$F \subseteq P[X]$*
obtains *g where $g \in P[X]$ and $g \neq 0$ and $\text{lcf } g = 1$ and *punit.reduced-GB F*
= {g} |
punit.reduced-GB F = {}
*<proof>**

corollary *deg-reduced-GB-univariate-le:*
assumes *$F \subseteq P[X]$ and $f \in \text{ideal } F$ and $f \neq 0$ and $g \in \text{punit.reduced-GB } F$*
shows *$\text{poly-deg } g \leq \text{poly-deg } f$*
<proof>

end

19.4 Homogeneity

lemma *is-reduced-GB-homogeneous*:

assumes $\bigwedge f. f \in F \implies \text{homogeneous } f$ **and** *punit.is-reduced-GB* G **and** *ideal* $G = \text{ideal } F$
and $g \in G$
shows *homogeneous* g
(*proof*)

lemma *lp-dehomogenize*:

assumes *is-hom-ord* x **and** *homogeneous* p
shows *lpp* (*dehomogenize* x p) = *except* (*lpp* p) $\{x\}$
(*proof*)

lemma *isGB-dehomogenize*:

assumes *is-hom-ord* x **and** *finite* X **and** $G \subseteq P[X]$ **and** *punit.is-Groebner-basis* G
and $\bigwedge g. g \in G \implies \text{homogeneous } g$
shows *punit.is-Groebner-basis* (*dehomogenize* x ‘ G)
(*proof*)

end

context *extended-ord-pm-powerprod*

begin

lemma *extended-ord-lp*:

assumes $\text{None} \notin \text{indets } p$
shows *restrict-indets-pp* (*extended-ord.lpp* p) = *lpp* (*restrict-indets* p)
(*proof*)

lemma *restrict-indets-reduced-GB*:

assumes *finite* X **and** $F \subseteq P[X]$
shows *punit.is-Groebner-basis* (*restrict-indets* ‘ *extended-ord.punit.reduced-GB* (*homogenize* None ‘ *extend-indets* ‘ F))
(*is* ?*thesis1*)
and *ideal* (*restrict-indets* ‘ *extended-ord.punit.reduced-GB* (*homogenize* None ‘ *extend-indets* ‘ F)) = *ideal* F
(*is* ?*thesis2*)
and *restrict-indets* ‘ *extended-ord.punit.reduced-GB* (*homogenize* None ‘ *extend-indets* ‘ F) $\subseteq P[X]$
(*is* ?*thesis3*)
(*proof*)

end

end

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