

An Isabelle/HOL formalisation of Green's Theorem

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Abstract

We formalise a statement of Greens theorem—the first formalisation to our knowledge—in Isabelle/HOL. The theorem statement that we formalise is enough for most applications, especially in physics and engineering. Our formalisation is made possible by a novel proof that avoids the ubiquitous line integral cancellation argument. This eliminates the need to formalise orientations and region boundaries explicitly with respect to the outwards-pointing normal vector. Instead we appeal to a homological argument about equivalences between paths.

1 Acknowledgements

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theory *General-Utils*

imports *HOL-Analysis.Analysis*

begin

lemma *lambda-skolem-gen*: $(\forall i. \exists f'::('a \wedge 'n) \Rightarrow 'a. P i f') \longleftrightarrow$
 $(\exists f'::('a \wedge 'n) \Rightarrow ('a \wedge 'n). \forall i. P i ((\lambda x. (f' x) \$ i)))$ (**is** *?lhs* \longleftrightarrow *?rhs*)

proof –

{ **assume** *H*: *?rhs*
then have *?lhs* **by** *auto* }

moreover

{ **assume** *H*: *?lhs*
then obtain *f''* **where** $f'':\forall i. P i (f'' i)$ **unfolding** *choice-iff*
by *metis*

let *?x* = $(\lambda x. (\chi i. (f'' i) x))$

{ **fix** *i*
from *f''* **have** $P i (f'' i)$ **by** *metis*
then have $P i (\lambda x. (?x x) \$ i)$ **by** *auto*

}
hence $\forall i. P i (\lambda x. (?x x) \$ i)$ **by** *metis*

hence *?rhs* **by** *metis*}

ultimately show *?thesis* by *metis*
qed

lemma *lambda-skolem-euclidean*: $(\forall i \in \text{Basis}. \exists f'::('a::\{\text{euclidean-space}\} \Rightarrow \text{real}). P\ i\ f') \longleftrightarrow$
 $(\exists f'::('a::\{\text{euclidean-space}\} \Rightarrow 'b::\{\text{euclidean-space}\}). \forall i \in \text{Basis}. P\ i\ ((\lambda x. (f'\ x) \cdot i)))$
(is *?lhs* \longleftrightarrow *?rhs*)

proof –

{ assume *H*: *?rhs*
then have *?lhs* by *auto* }
moreover
{ assume *H*: *?lhs*
then obtain *f''* where $f'':\forall i::('b::\{\text{euclidean-space}\}) \in \text{Basis}. P\ i\ (f''\ i)$ **unfolding**
choice-iff
by *metis*
let $?x = (\lambda x. (\sum i \in \text{Basis}. ((f''\ i)\ x) *_{\mathbb{R}} i))$
{ fix $i::'b::\{\text{euclidean-space}\}$
assume *ass*: $i \in \text{Basis}$
then have $P\ i\ (f''\ i)$
using *f''*
by *metis*
then have $P\ i\ (\lambda x. (?x\ x) \cdot i)$ **using** *ass* by *auto*
}
hence $*$: $\forall i \in \text{Basis}. P\ i\ (\lambda x. (?x\ x) \cdot i)$ by *auto*
then have *?rhs*
apply *auto*
proof
let $?f'6 = ?x$
show $\forall i \in \text{Basis}. P\ i\ (\lambda x. ?f'6\ x \cdot i)$ **using** $*$ by *auto*
qed }
ultimately show *?thesis* by *metis*
qed

lemma *lambda-skolem-euclidean-explicit*: $(\forall i \in \text{Basis}. \exists f'::('a::\{\text{euclidean-space}\} \Rightarrow \text{real}). P\ i\ f') \longleftrightarrow$
 $(\exists f'::('a::\{\text{euclidean-space}\} \Rightarrow 'a). \forall i \in \text{Basis}. P\ i\ ((\lambda x. (f'\ x) \cdot i)))$ (is *?lhs* \longleftrightarrow *?rhs*)

proof –

{ assume *H*: *?rhs*
then have *?lhs* by *auto* }
moreover
{ assume *H*: *?lhs*
then obtain *f''* where $f'':\forall i::('a::\{\text{euclidean-space}\}) \in \text{Basis}. P\ i\ (f''\ i)$ **unfolding**
choice-iff
by *metis*
let $?x = (\lambda x. (\sum i \in \text{Basis}. ((f''\ i)\ x) *_{\mathbb{R}} i))$
{ fix $i::'a::\{\text{euclidean-space}\}$

```

assume ass:  $i \in \text{Basis}$ 
then have  $P\ i\ (f''\ i)$ 
  using f''
  by metis
then have  $P\ i\ (\lambda x. (?x\ x) \cdot i)$  using ass by auto
}
hence  $*$ :  $\forall i \in \text{Basis}. P\ i\ (\lambda x. (?x\ x) \cdot i)$  by auto
then have ?rhs
  apply auto
proof
  let ?f'6 = ?x
  show  $\forall i \in \text{Basis}. P\ i\ (\lambda x. ?f'6\ x \cdot i)$  using  $*$  by auto
qed
ultimately show ?thesis by metis
qed

```

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lemma indic-ident:
 $\bigwedge (f::'a \Rightarrow \text{real})\ s. (\lambda x. (f\ x) * \text{indicator}\ s\ x) = (\lambda x. \text{if } x \in s \text{ then } f\ x \text{ else } 0)$ 
proof
  fix f::'a  $\Rightarrow$  real
  fix s::'a set
  fix x::'a
  show  $f\ x * \text{indicator}\ s\ x = (\text{if } x \in s \text{ then } f\ x \text{ else } 0)$ 
  by (simp add: indicator-def)
qed

```

```

lemma real-pair-basis:  $\text{Basis} = \{(1::\text{real}, 0::\text{real}), (0::\text{real}, 1::\text{real})\}$ 
by (simp add: Basis-prod-def insert-commute)

```

```

lemma real-singleton-in-borel:
shows  $\{a::\text{real}\} \in \text{sets borel}$ 
using Borel-Space.cbox-borel[of a a]
apply auto
done

```

```

lemma real-singleton-in-lborel:
shows  $\{a::\text{real}\} \in \text{sets lborel}$ 
using real-singleton-in-borel
apply auto
done

```

```

lemma cbox-diff:
shows  $\{0::\text{real}..1\} - \{0,1\} = \text{box } 0\ 1$ 
by (auto simp add: cbox-def)

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lemma sum-bij:

```

assumes *bij F*
 $\forall x \in s. f x = g (F x)$
shows $\bigwedge t. F^{-1} s = t \implies \text{sum } f s = \text{sum } g t$
by (*metis assms bij-betw-def bij-betw-subset subset-UNIV sum.reindex-cong*)

abbreviation *surj-on where*
 $\text{surj-on } s f \equiv s \subseteq \text{range } f$

lemma *surj-on-image-vimage-eq*: $\text{surj-on } s f \implies f^{-1} (f^{-1} s) = s$
by *fastforce*

end
theory *Derivs*
imports *General-Utills*
begin

lemma *field-simp-has-vector-derivative* [*derivative-intros*]:
 $(f \text{ has-field-derivative } y) F \implies (f \text{ has-vector-derivative } y) F$
by (*simp add: has-real-derivative-iff-has-vector-derivative*)

lemma *continuous-on-cases-empty* [*continuous-intros*]:
 $\llbracket \text{closed } S; \text{ continuous-on } S f; \bigwedge x. \llbracket x \in S; \neg P x \rrbracket \implies f x = g x \rrbracket \implies$
 $\text{continuous-on } S (\lambda x. \text{if } P x \text{ then } f x \text{ else } g x)$
using *continuous-on-cases [of - {}]* **by** *force*

lemma *inj-on-cases*:
assumes *inj-on f (Collect P \cap S) inj-on g (Collect (Not \circ P) \cap S)*
 $f^{-1} (\text{Collect } P \cap S) \cap g^{-1} (\text{Collect } (\text{Not } \circ P) \cap S) = \{\}$
shows *inj-on ($\lambda x. \text{if } P x \text{ then } f x \text{ else } g x$) S*
using *assms by (force simp: inj-on-def)*

lemma *inj-on-arccos*: $S \subseteq \{-1..1\} \implies \text{inj-on arccos } S$
by (*metis atLeastAtMost-iff cos-arccos inj-onI subsetCE*)

lemma *has-vector-derivative-componentwise-within*:
 $(f \text{ has-vector-derivative } f') \text{ (at } a \text{ within } S) \iff$
 $(\forall i \in \text{Basis}. ((\lambda x. f x \cdot i) \text{ has-vector-derivative } (f' \cdot i)) \text{ (at } a \text{ within } S))$
apply (*simp add: has-vector-derivative-def*)
apply (*subst has-derivative-componentwise-within*)
apply *simp*
done

lemma *has-vector-derivative-pair-within*:
fixes $f :: \text{real} \Rightarrow 'a::\text{euclidean-space}$ **and** $g :: \text{real} \Rightarrow 'b::\text{euclidean-space}$
assumes $\bigwedge u. u \in \text{Basis} \implies ((\lambda x. f x \cdot u) \text{ has-vector-derivative } f' \cdot u) \text{ (at } x \text{ within } S)$
 $\bigwedge u. u \in \text{Basis} \implies ((\lambda x. g x \cdot u) \text{ has-vector-derivative } g' \cdot u) \text{ (at } x \text{ within } S)$
shows $((\lambda x. (f x, g x)) \text{ has-vector-derivative } (f', g')) \text{ (at } x \text{ within } S)$
apply (*subst has-vector-derivative-componentwise-within*)

```

apply (auto simp: assms Basis-prod-def)
done

lemma piecewise-C1-differentiable-const:
  shows  $(\lambda x. c)$  piecewise-C1-differentiable-on S
  using continuous-on-const
  by (auto simp add: piecewise-C1-differentiable-on-def)

declare piecewise-C1-differentiable-const [simp, derivative-intros]
declare piecewise-C1-differentiable-neg [simp, derivative-intros]
declare piecewise-C1-differentiable-add [simp, derivative-intros]
declare piecewise-C1-differentiable-diff [simp, derivative-intros]

lemma piecewise-C1-differentiable-on-ident [simp, derivative-intros]:
  fixes  $f :: \text{real} \Rightarrow 'a::\text{real-normed-vector}$ 
  shows  $(\lambda x. x)$  piecewise-C1-differentiable-on S
  unfolding piecewise-C1-differentiable-on-def using C1-differentiable-on-ident
  by (blast intro: continuous-on-id C1-differentiable-on-ident)

lemma piecewise-C1-differentiable-on-mult [simp, derivative-intros]:
  fixes  $f :: \text{real} \Rightarrow 'a::\text{real-normed-algebra}$ 
  assumes  $f$  piecewise-C1-differentiable-on S  $g$  piecewise-C1-differentiable-on S
  shows  $(\lambda x. f x * g x)$  piecewise-C1-differentiable-on S
  using assms
  unfolding piecewise-C1-differentiable-on-def
  apply safe
  apply (blast intro: continuous-intros)
  apply (rename-tac A B)
  apply (rule-tac x=A  $\cup$  B in exI)
  apply (auto intro: C1-differentiable-on-mult C1-differentiable-on-subset)
  done

lemma C1-differentiable-on-cdiv [simp, derivative-intros]:
  fixes  $f :: \text{real} \Rightarrow 'a :: \text{real-normed-field}$ 
  shows  $f$  C1-differentiable-on S  $\implies (\lambda x. f x / c)$  C1-differentiable-on S
  by (simp add: divide-inverse)

lemma piecewise-C1-differentiable-on-cdiv [simp, derivative-intros]:
  fixes  $f :: \text{real} \Rightarrow 'a::\text{real-normed-field}$ 
  assumes  $f$  piecewise-C1-differentiable-on S
  shows  $(\lambda x. f x / c)$  piecewise-C1-differentiable-on S
  by (simp add: divide-inverse piecewise-C1-differentiable-const piecewise-C1-differentiable-on-mult
assms)

lemma sqrt-C1-differentiable [simp, derivative-intros]:
  assumes  $f: f$  C1-differentiable-on S and  $\text{fim}: f ' S \subseteq \{0 < ..\}$ 
  shows  $(\lambda x. \text{sqrt}(f x))$  C1-differentiable-on S

```

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proof –
  have contf: continuous-on S f
    by (simp add: C1-differentiable-imp-continuous-on f)
  show ?thesis
    using assms
    unfolding C1-differentiable-on-def has-real-derivative-iff-has-vector-derivative
    [symmetric]
    by (fastforce intro!: contf continuous-intros derivative-intros)
qed

lemma sqrt-piecewise-C1-differentiable [simp, derivative-intros]:
  assumes f: f piecewise-C1-differentiable-on S and fm: f ' S ⊆ {0<..}
  shows  $(\lambda x. \text{sqrt } (f x))$  piecewise-C1-differentiable-on S
  using assms
  unfolding piecewise-C1-differentiable-on-def
  by (fastforce intro!: continuous-intros derivative-intros)

lemma
  fixes f :: real  $\Rightarrow$  'a::{banach,real-normed-field}
  assumes f: f C1-differentiable-on S
  shows sin-C1-differentiable [simp, derivative-intros]:  $(\lambda x. \text{sin } (f x))$  C1-differentiable-on S
  and cos-C1-differentiable [simp, derivative-intros]:  $(\lambda x. \text{cos } (f x))$  C1-differentiable-on S
proof –
  have contf: continuous-on S f
    by (simp add: C1-differentiable-imp-continuous-on f)
  note df-sin = field-vector-diff-chain-at [where g=sin, unfolded o-def]
  note df-cos = field-vector-diff-chain-at [where g=cos, unfolded o-def]
  show  $(\lambda x. \text{sin } (f x))$  C1-differentiable-on S  $(\lambda x. \text{cos } (f x))$  C1-differentiable-on S
    using assms
    unfolding C1-differentiable-on-def has-real-derivative-iff-has-vector-derivative
    [symmetric]
    apply auto
    by (rule contf continuous-intros df-sin df-cos derivative-intros exI conjI ballI |
force)+
qed

lemma has-derivative-abs:
  fixes a::real
  assumes a  $\neq$  0
  shows  $(\text{abs has-derivative } ((*) (\text{sgn } a)))$  (at a)
proof –
  have [simp]: norm = abs
    using real-norm-def by force
  show ?thesis
    using has-derivative-norm [where 'a=real, simplified] assms
    by (simp add: mult-commute-abs)

```

qed

lemma *abs-C1-differentiable* [*simp*, *derivative-intros*]:

fixes $f :: \text{real} \Rightarrow \text{real}$

assumes $f: f \text{ C1-differentiable-on } S$ **and** $0 \notin f' S$

shows $(\lambda x. \text{abs } (f x)) \text{ C1-differentiable-on } S$

proof –

have *contf*: *continuous-on S f*

by (*simp add: C1-differentiable-imp-continuous-on f*)

note $df = \text{DERIV-chain}$ [**where** $f = \text{abs}$ **and** $g = f$, *unfolded o-def*]

show *?thesis*

using *assms*

unfolding *C1-differentiable-on-def has-real-derivative-iff-has-vector-derivative*
[*symmetric*]

apply *clarify*

apply (*rule df exI conjI ballI*)+

apply (*force simp: has-field-derivative-def intro: has-derivative-abs continuous-intros contf*)+

done

qed

lemma *C1-differentiable-on-pair* [*simp*, *derivative-intros*]:

fixes $f :: \text{real} \Rightarrow 'a::\text{euclidean-space}$ **and** $g :: \text{real} \Rightarrow 'b::\text{euclidean-space}$

assumes $f \text{ C1-differentiable-on } S$ $g \text{ C1-differentiable-on } S$

shows $(\lambda x. (f x, g x)) \text{ C1-differentiable-on } S$

using *assms unfolding C1-differentiable-on-def*

apply *safe*

apply (*rename-tac A B*)

apply (*intro exI ballI conjI*)

apply (*rule-tac f'=A x and g'=B x in has-vector-derivative-pair-within*)

using *has-vector-derivative-componentwise-within*

by (*blast intro: continuous-on-Pair*)+

lemma *piecewise-C1-differentiable-on-pair* [*simp*, *derivative-intros*]:

fixes $f :: \text{real} \Rightarrow 'a::\text{euclidean-space}$ **and** $g :: \text{real} \Rightarrow 'b::\text{euclidean-space}$

assumes $f \text{ piecewise-C1-differentiable-on } S$ $g \text{ piecewise-C1-differentiable-on } S$

shows $(\lambda x. (f x, g x)) \text{ piecewise-C1-differentiable-on } S$

using *assms unfolding piecewise-C1-differentiable-on-def*

by (*blast intro!: continuous-intros C1-differentiable-on-pair del: continuous-on-discrete intro: C1-differentiable-on-subset*)

lemma *test2*:

assumes $s: \bigwedge x. x \in \{0..1\} - s \implies g \text{ differentiable at } x$

and $fs: \text{finite } s$ **and** $uv: u \in \{0..1\} v \in \{0..1\} u \leq v$

and $x \in \{0..1\} x \notin (\lambda t. (v-u) *_R t + u) -' s$

shows $\text{vector-derivative } (\lambda x. g ((v-u) * x + u)) \text{ (at } x \text{ within } \{0..1\}) = (v-u)$
 $*_R \text{ vector-derivative } g \text{ (at } ((v-u) * x + u) \text{ within } \{0..1\})$

proof –

have $i:(g \text{ has-vector-derivative vector-derivative } g \text{ (at } ((v-u) * x + u))) \text{ (at$

$((v-u) * x + u)$
using *assms s [of (v - u) * x + u] uv mult-left-le [of x v-u]*
by *(auto simp: vector-derivative-works)*
have *ii: ((λx. g ((v - u) * x + u)) has-vector-derivative (v - u) *_R vector-derivative g (at ((v - u) * x + u))) (at x)*
by *(intro vector-diff-chain-at [simplified o-def] derivative-eq-intros | simp add: i)+*
have *0: 0 ≤ (v - u) * x + u*
using *assms uv by auto*
have *1: (v - u) * x + u ≤ 1*
using *assms uv*
by *simp (metis add.commute atLeastAtMost-iff atLeastatMost-empty-iff diff-ge-0-iff-ge empty-iff le-diff-eq mult-left-le)*
have *iii: vector-derivative g (at ((v - u) * x + u) within {0..1}) = vector-derivative g (at ((v - u) * x + u))*
using *Derivative.vector-derivative-at-within-ivl[OF i, of 0 1, OF 0 1]*
by *auto*
have *iv: vector-derivative (λx. g ((v - u) * x + u)) (at x within {0..1}) = (v - u) *_R vector-derivative g (at ((v - u) * x + u))*
using *Derivative.vector-derivative-at-within-ivl[OF ii, of 0 1] assms*
by *auto*
show *?thesis*
using *iii iv by auto*
qed

lemma *C1-differentiable-on-components:*

assumes $\bigwedge i. i \in \text{Basis} \implies (\lambda x. f x \cdot i)$ *C1-differentiable-on s*
shows *f C1-differentiable-on s*
proof *(clarsimp simp add: C1-differentiable-on-def has-vector-derivative-def)*
have **: $\forall f i x. x *_{\mathbb{R}} (f \cdot i) = (x *_{\mathbb{R}} f) \cdot i$ by auto*
have $\exists f'. \forall i \in \text{Basis}. \forall x \in s. ((\lambda x. f x \cdot i)$ *has-derivative* $(\lambda z. z *_{\mathbb{R}} f' x \cdot i))$ *(at x) \wedge continuous-on s f'*
using *assms lambda-skolem-euclidean[of λi D. (∀ x ∈ s. ((λx. f x · i) has-derivative (λz. z *_R D x)) (at x) \wedge continuous-on s D)*
apply *(simp only: C1-differentiable-on-def has-vector-derivative-def *)*
using *continuous-on-componentwise[of s]*
by *metis*
then obtain *f' where f': $\forall i \in \text{Basis}. \forall x \in s. ((\lambda x. f x \cdot i)$ has-derivative* $(\lambda z. z *_{\mathbb{R}} f' x \cdot i))$ *(at x) \wedge continuous-on s f'*
by *auto*
then have *0: (∀ x ∈ s. (f has-derivative (λz. z *_R f' x)) (at x)) \wedge continuous-on s f'*
using *f' has-derivative-componentwise-within[of f, where S= UNIV]*
by *auto*
then show $\exists D. (\forall x \in s. (f$ *has-derivative* $(\lambda z. z *_{\mathbb{R}} D x))$ *(at x)) \wedge continuous-on s D* **by** *metis*
qed

lemma *piecewise-C1-differentiable-on-components:*

assumes *finite t*
 $\bigwedge i. i \in \text{Basis} \implies (\lambda x. f x \cdot i) \text{ C1-differentiable-on } s - t$
 $\bigwedge i. i \in \text{Basis} \implies \text{continuous-on } s (\lambda x. f x \cdot i)$
shows *f piecewise-C1-differentiable-on s*
using *C1-differentiable-on-components assms continuous-on-componentwise piecewise-C1-differentiable-on-def* **by** *blast*

lemma *all-components-smooth-one-pw-smooth-is-pw-smooth:*

assumes $\bigwedge i. i \in \text{Basis} - \{j\} \implies (\lambda x. f x \cdot i) \text{ C1-differentiable-on } s$
assumes $(\lambda x. f x \cdot j) \text{ piecewise-C1-differentiable-on } s$
shows *f piecewise-C1-differentiable-on s*
proof –
have *is-cont: $\forall i \in \text{Basis}. \text{continuous-on } s (\lambda x. f x \cdot i)$*
using *assms C1-differentiable-imp-continuous-on piecewise-C1-differentiable-on-def*
by *fastforce*
obtain *t where t:(finite t $\wedge (\lambda x. f x \cdot j) \text{ C1-differentiable-on } s - t)$* **using**
assms(2) piecewise-C1-differentiable-on-def **by** *auto*
show *?thesis*
using *piecewise-C1-differentiable-on-components[where ?f = f]*
using *assms(2) piecewise-C1-differentiable-on-def*
 $\text{C1-differentiable-on-subset}[OF \text{ assms(1) Diff-subset, where ?B1 = t}] \text{ t is-cont}$
by *fastforce*
qed

lemma *derivative-component-fun-component:*

fixes *i::'a::euclidean-space*
assumes *f differentiable (at x)*
shows $((\text{vector-derivative } f \text{ (at } x)) \cdot i) = ((\text{vector-derivative } (\lambda x. (f x) \cdot i) \text{ (at } x)))$
proof –
have $(\lambda x. f x \cdot i) \text{ has-vector-derivative } \text{vector-derivative } f \text{ (at } x) \cdot i \text{ (at } x)$
using *assms and bounded-linear.has-vector-derivative[of $(\lambda x. x \cdot i) f$ (vector-derivative f (at x)) (at x)]* **and**
 $\text{bounded-linear-inner-left}[of i]$ **and** $\text{vector-derivative-works}[of f \text{ (at } x)]$
by *blast*
then show $((\text{vector-derivative } f \text{ (at } x)) \cdot i) = ((\text{vector-derivative } (\lambda x. (f x) \cdot i) \text{ (at } x)))$
using $\text{vector-derivative-works}[of (\lambda x. (f x) \cdot i) \text{ (at } x)]$ **and**
 $\text{differentiableI-vector}[of (\lambda x. (f x) \cdot i) \text{ (vector-derivative } f \text{ (at } x) \cdot i) \text{ (at } x)]$
and
 $\text{Derivative.vector-derivative-at}$
by *force*
qed

lemma *gamma-deriv-at-within:*

assumes *a-leq-b: $a < b$* **and**
 $x\text{-within-bounds: } x \in \{a..b\}$ **and**
 $\text{gamma-differentiable: } \forall x \in \{a..b\}. \gamma \text{ differentiable at } x$
shows $\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{a..b\}) = \text{vector-derivative } \gamma \text{ (at } x)$

using *Derivative.vector-derivative-at-within-ivl*[of γ vector-derivative γ (at x) x
 a b]
gamma-differentiable x -within-bounds a -leq- b
by (*auto simp add: vector-derivative-works*)

lemma *islimpt-diff-finite*:
assumes *finite* ($t::'a::t1$ -space set)
shows x *islimpt* $s - t = x$ *islimpt* s
proof –
have *iii*: $s - t = s - (t \cap s)$ **by** *auto*
have $(t \cap s) \subseteq s$ **by** *auto*
have *ii*: *finite* $(t \cap s)$ **using** *assms(1)* **by** *auto*
have *i*: $(t \cap s) \cup (s - (t \cap s)) = (s)$
using *assms* **by** *auto*
then have x *islimpt* $s - (t \cap s) = x$ *islimpt* s
by (*metis ii islimpt-Un-finite*)
then show *?thesis* **using** *iii* **by** *auto*
qed

lemma *ivl-limpt-diff*:
assumes *finite* s $a < b$ ($x::real$) $\in \{a..b\} - s$
shows x *islimpt* $\{a..b\} - s$
proof –
have x *islimpt* $\{a..b\}$
proof (*cases* $x \in \{a,b\}$)
have *i*: *finite* $\{a,b\}$ **and** *ii*: $\{a, b\} \cup \{a <..<b\} = \{a..b\}$ **using** *assms* **by** *auto*
assume $x \in \{a,b\}$
then show *?thesis*
by (*meson DiffE assms(2) assms(3) islimpt-Icc*)
next
assume $x \notin \{a,b\}$
then show x *islimpt* $\{a..b\}$ **using** *assms* **by** *auto*
qed
then show x *islimpt* $\{a..b\} - s$ **using** *islimpt-diff-finite[OF assms(1)] assms*
by *fastforce*
qed

lemma *ivl-closure-diff-del*:
assumes *finite* s $a < b$ ($x::real$) $\in \{a..b\} - s$
shows $x \in$ *closure* ($(\{a..b\} - s) - \{x\}$)
using *ivl-limpt-diff islimpt-in-closure assms* **by** *blast*

lemma *ivl-not-trivial-limit-within*:
assumes *finite* s
 $a < b$
 $(x::real) \in \{a..b\} - s$
shows *at* x *within* $\{a..b\} - s \neq$ *bot*
using *assms ivl-closure-diff-del not-trivial-limit-within*
by *blast*

lemma *vector-derivative-at-within-non-trivial-limit:*
at x within s $\neq \text{bot}$ \wedge (*f has-vector-derivative f'*) (*at x*) \implies
vector-derivative f (*at x within s*) = *f'*
using *has-vector-derivative-at-within* *vector-derivative-within* **by** *fastforce*

lemma *vector-derivative-at-within-ivl-diff:*
finite s \wedge *a < b* \wedge (*x::real*) \in {*a..b*} - *s* \wedge (*f has-vector-derivative f'*) (*at x*) \implies
vector-derivative f (*at x within {a..b} - s*) = *f'*
using *vector-derivative-at-within-non-trivial-limit* *ivl-not-trivial-limit-within* **by**
fastforce

lemma *gamma-deriv-at-within-diff:*
assumes *a-leq-b: a < b* **and**
x-within-bounds: x \in {*a..b*} - *s* **and**
gamma-differentiable: $\forall x \in \{a .. b\} - s.$ γ *differentiable at x* **and**
s-subset: s \subseteq {*a..b*} **and**
finite-s: finite s
shows *vector-derivative γ* (*at x within {a..b} - s*)
= *vector-derivative γ* (*at x*)
using *vector-derivative-at-within-ivl-diff* [*of s a b x γ vector-derivative γ* (*at x*)]
gamma-differentiable
x-within-bounds a-leq-b s-subset finite-s
by (*auto simp add: vector-derivative-works*)

lemma *gamma-deriv-at-within-gen:*
assumes *a-leq-b: a < b* **and**
x-within-bounds: x \in *s* **and**
s-subset: s \subseteq {*a..b*} **and**
gamma-differentiable: $\forall x \in s.$ γ *differentiable at x*
shows *vector-derivative γ* (*at x within ({a..b})*) = *vector-derivative γ* (*at x*)
using *Derivative.vector-derivative-at-within-ivl*[*of γ vector-derivative γ* (*at x*) *x*
a b]
gamma-differentiable x-within-bounds a-leq-b s-subset
by (*auto simp add: vector-derivative-works*)

lemma *derivative-component-fun-component-at-within-gen:*
assumes *gamma-differentiable: $\forall x \in s.$* γ *differentiable at x* **and** *s-subset: s* \subseteq
{*0..1*}
shows $\forall x \in s.$ *vector-derivative* ($\lambda x. \gamma x$) (*at x within {0..1}*) \cdot (*i::'a:: eu-*
clidean-space)
= *vector-derivative* ($\lambda x. \gamma x \cdot i$) (*at x within {0..1}*)

proof -

have *gamma-i-component-smooth:*

$\forall x \in s.$ ($\lambda x. \gamma x \cdot i$) *differentiable at x*

using *gamma-differentiable*

by *auto*

show $\forall x \in s.$ *vector-derivative* ($\lambda x. \gamma x$) (*at x within {0..1}*) \cdot *i*
= *vector-derivative* ($\lambda x. \gamma x \cdot i$) (*at x within {0..1}*)

```

proof
  fix  $x::\text{real}$ 
  assume  $x\text{-within-bounds}: x \in s$ 
  have  $\text{gamma-deriv-at-within}$ :
     $\text{vector-derivative } (\lambda x. \gamma x) \text{ (at } x \text{ within } \{0..1\}) = \text{vector-derivative } (\lambda x. \gamma$ 
 $x) \text{ (at } x)$ 
    using  $\text{gamma-deriv-at-within-gen[of } 0 \ 1] \ x\text{-within-bounds}$ 
     $\text{gamma-differentiable } s\text{-subset}$ 
    by  $(\text{auto simp add: vector-derivative-works})$ 
  then have  $\text{gamma-component-deriv-at-within}$ :
     $\text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x)$ 
     $= \text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x \text{ within } \{0..1\})$ 
  using  $\text{gamma-deriv-at-within-gen[of } 0 \ 1, \text{ where } \gamma = (\lambda x. \gamma x \cdot i)] \ x\text{-within-bounds}$ 
     $\text{gamma-i-component-smooth } s\text{-subset}$ 
    by  $(\text{auto simp add: vector-derivative-works})$ 
  have  $\text{gamma-component-deriv-eq-gamma-deriv-component}$ :
     $\text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x) = \text{vector-derivative } (\lambda x. \gamma x) \text{ (at } x) \cdot i$ 
    using  $\text{derivative-component-fun-component[of } \gamma \ x \ i] \ \text{gamma-differentiable}$ 
 $x\text{-within-bounds}$ 
    by  $\text{auto}$ 
  show  $\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i = \text{vector-derivative } (\lambda x. \gamma x \cdot$ 
 $i) \text{ (at } x \text{ within } \{0..1\})$ 
    using  $\text{gamma-component-deriv-eq-gamma-deriv-component } \text{gamma-component-deriv-at-within}$ 
 $\text{gamma-deriv-at-within}$ 
    by  $\text{auto}$ 
  qed
qed

```

```

lemma  $\text{derivative-component-fun-component-at-within}$ :
  assumes  $\text{gamma-differentiable}: \forall x \in \{0 .. 1\}. \gamma \text{ differentiable at } x$ 
  shows  $\forall x \in \{0..1\}. \text{vector-derivative } (\lambda x. \gamma x) \text{ (at } x \text{ within } \{0..1\}) \cdot (i::'a::$ 
 $\text{euclidean-space})$ 
     $= \text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x \text{ within } \{0..1\})$ 

```

proof –

```

have  $\text{gamma-i-component-smooth}$ :
   $\forall x \in \{0 .. 1\}. (\lambda x. \gamma x \cdot i) \text{ differentiable at } x$ 
  using  $\text{gamma-differentiable}$  by  $\text{auto}$ 
show  $\forall x \in \{0..1\}. \text{vector-derivative } (\lambda x. \gamma x) \text{ (at } x \text{ within } \{0..1\}) \cdot i$ 
     $= \text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x \text{ within } \{0..1\})$ 

```

proof

```

fix  $x::\text{real}$ 
assume  $x\text{-within-bounds}: x \in \{0..1\}$ 
have  $\text{gamma-deriv-at-within}$ :
   $\text{vector-derivative } (\lambda x. \gamma x) \text{ (at } x \text{ within } \{0..1\}) = \text{vector-derivative } (\lambda x. \gamma x)$ 
 $(\text{at } x)$ 
  using  $\text{gamma-deriv-at-within[of } 0 \ 1] \ x\text{-within-bounds}$ 
   $\text{gamma-differentiable}$ 
  by  $(\text{auto simp add: vector-derivative-works})$ 
have  $\text{gamma-component-deriv-at-within}$ :

```

```

      vector-derivative (λx. γ x · i) (at x) = vector-derivative (λx. γ x · i) (at x
within {0..1})
    using Derivative.vector-derivative-at-within-ivl[of (λx. (γ x) · i) vector-derivative
(λx. (γ x) · i) (at x) x 0 1]
      has-vector-derivative-at-within[of (λx. γ x · i) vector-derivative (λx. γ x ·
i) (at x) x {0..1}]
      gamma-i-component-smooth x-within-bounds
    by (simp add: vector-derivative-works)
  have gamma-component-deriv-eq-gamma-deriv-component:
    vector-derivative (λx. γ x · i) (at x) = vector-derivative (λx. γ x) (at x) · i
    using derivative-component-fun-component[of γ x i] gamma-differentiable
x-within-bounds
    by auto
  show vector-derivative γ (at x within {0..1}) · i = vector-derivative (λx. γ x ·
i) (at x within {0..1})
    using gamma-component-deriv-eq-gamma-deriv-component gamma-component-deriv-at-within
gamma-deriv-at-within
    by auto
  qed
qed

```

```

lemma straight-path-differentiable-x:
  fixes b :: real and y1 :: real
  assumes gamma-def: γ = (λx. (b, y2 + y1 * x))
  shows ∀x. γ differentiable at x
  unfolding gamma-def differentiable-def
  by (fast intro!: derivative-intros)

```

```

lemma straight-path-differentiable-y:
  fixes b :: real and
  y1 y2 :: real
  assumes gamma-def: γ = (λx. (y2 + y1 * x, b))
  shows ∀x. γ differentiable at x
  unfolding gamma-def differentiable-def
  by (fast intro!: derivative-intros)

```

```

lemma piecewise-C1-differentiable-on-imp-continuous-on:
  assumes f piecewise-C1-differentiable-on s
  shows continuous-on s f
  using assms
  by (auto simp add: piecewise-C1-differentiable-on-def)

```

```

lemma boring-lemma1:
  fixes f :: real ⇒ real
  assumes (f has-vector-derivative D) (at x)
  shows ((λx. (f x, 0)) has-vector-derivative ((D, 0 :: real))) (at x)
proof -
  have *: ((λx. (f x) *R (1, 0)) has-vector-derivative (D *R (1, 0))) (at x)
  using bounded-linear.has-vector-derivative[OF Real-Vector-Spaces.bounded-linear-scaleR-left

```

```

assms(1),
  of (1,0)] by auto
have ((λx. (f x) *R (1,0)) has-vector-derivative (D,0)) (at x)
proof -
  have (D, 0::'a) = D *R (1, 0)
  by simp
  then show ?thesis
  by (metis (no-types) *)
qed
then show ?thesis by auto
qed

lemma boring-lemma2:
  fixes f :: real ⇒ real
  assumes (f has-vector-derivative D) (at x)
  shows ((λx. (0, f x)) has-vector-derivative (0, D)) (at x)
proof -
  have *: ((λx. (f x) *R (0,1)) has-vector-derivative (D *R (0,1))) (at x)
  using bounded-linear.has-vector-derivative[OF Real-Vector-Spaces.bounded-linear-scaleR-left
assms(1),
  of (0,1)] by auto
  then have ((λx. (f x) *R (0,1)) has-vector-derivative ((0,D))) (at x)
  using scaleR-Pair Real-Vector-Spaces.real-scaleR-def
proof -
  have (0::'b, D) = D *R (0, 1)
  by auto
  then show ?thesis
  by (metis (no-types) *)
qed
then show ?thesis by auto
qed

lemma pair-prod-smooth-pw-smooth:
  assumes (f::real ⇒ real) C1-differentiable-on s (g::real ⇒ real) piecewise-C1-differentiable-on
  s
  shows (λx. (f x, g x)) piecewise-C1-differentiable-on s
proof -
  have f-cont: continuous-on s f
  using assms(1) C1-differentiable-imp-continuous-on
  by fastforce
  have g-cont: continuous-on s g
  using assms(2) by (auto simp add: piecewise-C1-differentiable-on-def)
  obtain t where t:(finite t ∧ g C1-differentiable-on s - t) using assms(2) piece-
wise-C1-differentiable-on-def by auto
  show ?thesis
  using piecewise-C1-differentiable-on-components[where ?f = (λx. (f x, g x))]
  apply (simp add: real-pair-basis)
  using assms(2) piecewise-C1-differentiable-on-def
  C1-differentiable-on-subset[OF assms(1) Diff-subset, where ?B1 = t] t

```

f-cont g-cont
 by *fastforce*
 qed

lemma *scale-shift-smooth*:
 shows $(\lambda x. a + b * x)$ *C1-differentiable-on s*
proof –
 show $(\lambda x. a + b * x)$ *C1-differentiable-on s*
 using *C1-differentiable-on-mult C1-differentiable-on-add C1-differentiable-on-const*
C1-differentiable-on-ident by *auto*
 qed

lemma *open-diff*:
 assumes *finite* ($t :: 'a :: t1\text{-space}$ set)
 open ($s :: 'a$ set)
 shows *open* ($s - t$)
 using *assms*
proof(*induction t*)
 show *open s* \implies *open* ($s - \{\}$) by *auto*
next
 fix $x :: 'a :: t1\text{-space}$
 fix $F :: 'a :: t1\text{-space}$ set
 assume *step*: *finite F* $x \notin F$ *open s*
 then have *i*: $(s - \text{insert } x F) = (s - F) - \{x\}$ by *auto*
 assume *ind-hyp*: (*open s* \implies *open* ($s - F$))
 show *open* ($s - \text{insert } x F$)
 apply (*simp only: i*)
 using *open-delete[of s - F]* *ind-hyp[OF step(3)]* by *auto*
 qed

lemma *has-derivative-transform-within*:
 assumes $0 < d$
 and $x \in s$
 and $\forall x' \in s. \text{dist } x' x < d \implies f x' = g x'$
 and (*f has-derivative f'*) (*at x within s*)
 shows (*g has-derivative f'*) (*at x within s*)
 using *assms*
unfolding *has-derivative-within*
 by (*force simp add: intro: Lim-transform-within*)

lemma *has-derivative-transform-within-ivl*:
 assumes $(0 :: \text{real}) < b$
 and $\forall x \in \{a..b\} - s. f x = g x$
 and $x \in \{a..b\} - s$
 and (*f has-derivative f'*) (*at x within* $\{a..b\} - s$)
 shows (*g has-derivative f'*) (*at x within* $\{a..b\} - s$)
 using *has-derivative-transform-within[of b x {a..b} - s]* *assms*
 by *auto*

lemma *has-vector-derivative-transform-within-ivl*:
assumes $(0::\text{real}) < b$
and $\forall x \in \{a..b\} - s. f\ x = g\ x$
and $x \in \{a..b\} - s$
and $(f\ \text{has-vector-derivative}\ f')$ (at x within $\{a..b\} - s$)
shows $(g\ \text{has-vector-derivative}\ f')$ (at x within $\{a..b\} - s$)
using *assms has-derivative-transform-within-ivl*
apply (*auto simp add: has-vector-derivative-def*)
by *blast*

lemma *has-derivative-transform-at*:
assumes $0 < d$
and $\forall x'. \text{dist}\ x'\ x < d \longrightarrow f\ x' = g\ x'$
and $(f\ \text{has-derivative}\ f')$ (at x)
shows $(g\ \text{has-derivative}\ f')$ (at x)
using *has-derivative-transform-within [of d x UNIV f g f'] assms*
by *simp*

lemma *has-vector-derivative-transform-at*:
assumes $0 < d$
and $\forall x'. \text{dist}\ x'\ x < d \longrightarrow f\ x' = g\ x'$
and $(f\ \text{has-vector-derivative}\ f')$ (at x)
shows $(g\ \text{has-vector-derivative}\ f')$ (at x)
using *assms*
unfolding *has-vector-derivative-def*
by (*rule has-derivative-transform-at*)

lemma *C1-diff-components-2*:
assumes $b \in \text{Basis}$
assumes $f\ \text{C1-differentiable-on}\ s$
shows $(\lambda x. f\ x \cdot b)\ \text{C1-differentiable-on}\ s$
proof –
obtain D **where** $D: (\forall x \in s. (f\ \text{has-derivative}\ (\lambda z. z *_{\mathbb{R}} D\ x))\ (at\ x))\ \text{continuous-on}\ s\ D$
using *assms(2) by (fastforce simp add: C1-differentiable-on-def has-vector-derivative-def)*
show *?thesis*
proof (*simp add: C1-differentiable-on-def has-vector-derivative-def, intro exI conjI*)
show *continuous-on s* $(\lambda x. D\ x \cdot b)$ **using** D *continuous-on-componentwise*
assms(1) by fastforce
show $(\forall x \in s. ((\lambda x. f\ x \cdot b)\ \text{has-derivative}\ (\lambda y. y * (\lambda x. D\ x \cdot b)\ x))\ (at\ x))$
using *has-derivative-inner-left D(1) by fastforce*
qed
qed

lemma *eq-smooth*:
assumes $0 < d$
 $\forall x \in s. \forall y. \text{dist}\ x\ y < d \longrightarrow f\ y = g\ y$
 $f\ \text{C1-differentiable-on}\ s$

shows g *C1-differentiable-on* s
proof (*simp add: C1-differentiable-on-def*)
obtain D **where** D : $(\forall x \in s. (f \text{ has-vector-derivative } D \ x) \ (at \ x)) \wedge \text{continuous-on } s \ D$
using *assms* **by** (*auto simp add: C1-differentiable-on-def*)
then have f : $(\forall x \in s. (g \text{ has-vector-derivative } D \ x) \ (at \ x))$
using *assms(1-2)*
by (*metis dist-commute has-vector-derivative-transform-at*)
have $(\forall x \in s. (g \text{ has-vector-derivative } D \ x) \ (at \ x)) \wedge \text{continuous-on } s \ D$ **using** D
by *auto*
then show $\exists D. (\forall x \in s. (g \text{ has-vector-derivative } D \ x) \ (at \ x)) \wedge \text{continuous-on } s \ D$ **by** *metis*
qed

lemma *eq-pw-smooth*:

assumes $0 < d$
 $\forall x \in s. \forall y. \text{dist } x \ y < d \longrightarrow f \ y = g \ y$
 $\forall x \in s. f \ x = g \ x$
 f *piecewise-C1-differentiable-on* s
shows g *piecewise-C1-differentiable-on* s
proof (*simp add: piecewise-C1-differentiable-on-def*)
have g -*cont*: *continuous-on* $s \ g$ **using** *assms piecewise-C1-differentiable-const*
by (*simp add: piecewise-C1-differentiable-on-def*)
obtain t **where** t : *finite* $t \wedge f$ *C1-differentiable-on* $s - t$
using *assms* **by** (*auto simp add: piecewise-C1-differentiable-on-def*)
then have g *C1-differentiable-on* $s - t$ **using** *assms eq-smooth* **by** *blast*
then show *continuous-on* $s \ g \wedge (\exists t. \text{finite } t \wedge g \text{ C1-differentiable-on } s - t)$
using t g -*cont* **by** *metis*
qed

lemma *scale-piecewise-C1-differentiable-on*:

assumes f *piecewise-C1-differentiable-on* s
shows $(\lambda x. (c::\text{real}) * (f \ x))$ *piecewise-C1-differentiable-on* s
proof (*simp add: piecewise-C1-differentiable-on-def, intro conjI*)
show *continuous-on* $s \ (\lambda x. c * f \ x)$
using *assms continuous-on-mult-left*
by (*auto simp add: piecewise-C1-differentiable-on-def*)
show $\exists t. \text{finite } t \wedge (\lambda x. c * f \ x)$ *C1-differentiable-on* $s - t$
using *assms continuous-on-mult-left*
by (*auto simp add: piecewise-C1-differentiable-on-def*)
qed

lemma *eq-smooth-gen*:

assumes f *C1-differentiable-on* s
 $\forall x. f \ x = g \ x$
shows g *C1-differentiable-on* s
using *assms* **unfolding** *C1-differentiable-on-def*
by (*metis (no-types, lifting) has-vector-derivative-weaken UNIV-I top-greatest*)

```

lemma subpath-compose:
  shows (subpath a b  $\gamma$ ) =  $\gamma$  o ( $\lambda x. (b - a) * x + a$ )
  by (auto simp add: subpath-def)

lemma subpath-smooth:
  assumes  $\gamma$  C1-differentiable-on {0..1}  $0 \leq a < b \leq 1$ 
  shows (subpath a b  $\gamma$ ) C1-differentiable-on {0..1}
proof -
  have  $\gamma$  C1-differentiable-on {a..b}
    apply (rule C1-differentiable-on-subset)
    using assms by auto
  then have  $\gamma$  C1-differentiable-on ( $\lambda x. (b - a) * x + a$ ) ‘ {0..1}
    using  $\langle a < b \rangle$  closed-segment-eq-real-ivl closed-segment-real-eq by auto
  moreover have finite ({0..1}  $\cap$  ( $\lambda x. (b - a) * x + a$ ) - ‘ {x}) for x
proof -
  have (( $\lambda x. (b - a) * x + a$ ) - ‘ {x}) = {(x - a) / (b - a)}
    using assms by (auto simp add: divide-simps)
  then show ?thesis
    by auto
qed
  ultimately show ?thesis
    by (force simp add: subpath-compose intro: C1-differentiable-compose derivative-intros)
qed

lemma has-vector-derivative-divide[derivative-intros]:
  fixes a :: 'a::real-normed-field'
  shows (f has-vector-derivative x) F  $\implies$  (( $\lambda x. f\ x / a$ ) has-vector-derivative (x / a)) F
  unfolding divide-inverse by (fact has-vector-derivative-mult-left)

end
theory Integrals
  imports HOL-Analysis.Analysis General-Utils
begin

lemma gauge-integral-Fubini-universe-x:
  fixes f :: ('a::euclidean-space * 'b::euclidean-space)  $\Rightarrow$  'c::euclidean-space'
  assumes fun-lesbegue-integrable: integrable lborel f and
    x-axis-integral-measurable: ( $\lambda x. \text{integral UNIV } (\lambda y. f(x, y))$ )  $\in$  borel-measurable lborel
  shows integral UNIV f = integral UNIV ( $\lambda x. \text{integral UNIV } (\lambda y. f(x, y))$ )
    ( $\lambda x. \text{integral UNIV } (\lambda y. f(x, y))$ ) integrable-on UNIV
proof -
  have f-is-measurable: f  $\in$  borel-measurable lborel
    using fun-lesbegue-integrable and borel-measurable-integrable
    by auto
  have fun-lborel-prod-integrable:
    integrable (lborel  $\otimes_M$  lborel) f

```

using *fun-lesbegue-integrable*
by (*simp add: lborel-prod*)
then have *region-integral-is-one-twoD-integral*:
 $(LBINT x. LBINT y. f(x, y)) = integral^L (lborel \otimes_M lborel) f$
using *lborel-pair.integral-fst'*
by *auto*
then have *AE-one-D-integrals-eq*: $AE x$ in *lborel*. $(LBINT y. f(x, y)) = integral$
 $UNIV (\lambda y. f(x, y))$
proof –
have $AE x$ in *lborel*. *integrable lborel* $(\lambda y. f(x, y))$
using *lborel-pair.AE-integrable-fst'* **and** *fun-lborel-prod-integrable*
by *blast*
then show *?thesis*
using *integral-lborel* **and** *always-eventually*
and *AE-mp*
by *fastforce*
qed
have *one-D-integral-measurable*:
 $(\lambda x. LBINT y. f(x, y)) \in borel-measurable lborel$
using *f-is-measurable* **and** *lborel.borel-measurable-lebesgue-integral*
by *auto*
then have *second-lesbegue-integral-eq*:
 $(LBINT x. LBINT y. f(x, y)) = (LBINT x. integral UNIV (\lambda y. f(x, y)))$
using *x-axis-integral-measurable* **and** *integral-cong-AE* **and** *AE-one-D-integrals-eq*
by *blast*
have *integrable lborel* $(\lambda x. LBINT y. f(x, y))$
using *fun-lborel-prod-integrable* **and** *lborel-pair.integrable-fst'*
by *auto*
then have *oneD-gauge-integral-lesbegue-integrable*:
integrable lborel $(\lambda x. integral UNIV (\lambda y. f(x, y)))$
using *x-axis-integral-measurable* **and** *AE-one-D-integrals-eq* **and** *integrable-cong-AE-imp*
by *blast*
then show *one-D-gauge-integral-integrable*:
 $(\lambda x. integral UNIV (\lambda y. f(x, y)))$ *integrable-on UNIV*
using *integrable-on-lborel*
by *auto*
have $(LBINT x. integral UNIV (\lambda y. f(x, y))) = integral UNIV (\lambda x. integral UNIV$
 $(\lambda y. f(x, y)))$
using *integral-lborel oneD-gauge-integral-lesbegue-integrable*
by *fastforce*
then have *twoD-lesbeuge-eq-twoD-gauge*:
 $(LBINT x. LBINT y. f(x, y)) = integral UNIV (\lambda x. integral UNIV (\lambda y. f(x,$
 $y)))$
using *second-lesbegue-integral-eq*
by *auto*
then show *integral UNIV f = integral UNIV* $(\lambda x. integral UNIV (\lambda y. f(x, y)))$
using *fun-lesbegue-integrable* **and** *integral-lborel* **and** *region-integral-is-one-twoD-integral*
by (*metis lborel-prod*)
qed

lemma *gauge-integral-Fubini-universe-y*:
fixes $f :: ('a::\text{euclidean-space} * 'b::\text{euclidean-space}) \Rightarrow 'c::\text{euclidean-space}$
assumes *fun-lesbegue-integrable*: *integrable lborel f* **and**
y-axis-integral-measurable: $(\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) \in \text{borel-measurable lborel}$
shows $\text{integral UNIV } f = \text{integral UNIV } (\lambda x. \text{integral UNIV } (\lambda y. f(y, x)))$
 $(\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) \text{ integrable-on UNIV}$
proof –
have *f-is-measurable*: $f \in \text{borel-measurable lborel}$
using *fun-lesbegue-integrable* **and** *borel-measurable-integrable*
by *auto*
have *fun-lborel-prod-integrable*:
integrable (lborel \otimes_M lborel) f
using *fun-lesbegue-integrable*
by (*simp add: lborel-prod*)
then have *region-integral-is-one-twoD-integral*:
 $(\text{LBINT } x. \text{LBINT } y. f(y, x)) = \text{integral}^L (\text{lborel} \otimes_M \text{lborel}) f$
by (*simp add: lborel-pair.integrable-product-swap-iff lborel-pair.integral-fst lborel-pair.integral-product-swap*)
then have *AE-one-D-integrals-eq*: $AE \ x \text{ in } \text{lborel}. (\text{LBINT } y. f(y, x)) = \text{integral UNIV } (\lambda y. f(y, x))$
proof –
have $AE \ x \text{ in } \text{lborel}. \text{integrable lborel } (\lambda y. f(y, x))$
using *lborel-pair.AE-integrable-fst'* **and** *fun-lborel-prod-integrable*
lborel-pair.AE-integrable-fst lborel-pair.integrable-product-swap
by *blast*
then show *?thesis*
using *integral-lborel always-eventually AE-mp* **by** *fastforce*
qed
have *one-D-integral-measurable*:
 $(\lambda x. \text{LBINT } y. f(y, x)) \in \text{borel-measurable lborel}$
using *f-is-measurable* **and** *lborel.borel-measurable-lebesgue-integral*
by *auto*
then have *second-lesbegue-integral-eq*:
 $(\text{LBINT } x. \text{LBINT } y. f(y, x)) = (\text{LBINT } x. \text{integral UNIV } (\lambda y. f(y, x)))$
using *y-axis-integral-measurable* **and** *integral-cong-AE* **and** *AE-one-D-integrals-eq*
by *blast*
have *integrable lborel* $(\lambda x. \text{LBINT } y. f(y, x))$
using *fun-lborel-prod-integrable* **and** *lborel-pair.integrable-fst'*
by (*simp add: lborel-pair.integrable-fst lborel-pair.integrable-product-swap*)
then have *oneD-gauge-integral-lesbegue-integrable*:
integrable lborel $(\lambda x. \text{integral UNIV } (\lambda y. f(y, x)))$
using *y-axis-integral-measurable* **and** *AE-one-D-integrals-eq* **and** *integrable-cong-AE-imp*
by *blast*
then show *one-D-gauge-integral-integrable*:
 $(\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) \text{ integrable-on UNIV}$
using *integrable-on-lborel* **by** *auto*
have $(\text{LBINT } x. \text{integral UNIV } (\lambda y. f(y, x))) = \text{integral UNIV } (\lambda x. \text{integral UNIV } (\lambda y. f(y, x)))$

using *integral-lborel oneD-gauge-integral-lesbegue-integrable*
by *fastforce*
then have *twoD-lesbeuge-eq-twoD-gauge*:
 $(LBINT\ x.\ LBINT\ y.\ f\ (y,\ x)) = integral\ UNIV\ (\lambda x.\ integral\ UNIV\ (\lambda y.\ f(y,\ x)))$
using *second-lesbegue-integral-eq* **by** *auto*
then show *integral UNIV f = integral UNIV (\lambda x. integral UNIV (\lambda y. f(y, x)))*
using *fun-lesbegue-integrable* **and** *integral-lborel* **and** *region-integral-is-one-twoD-integral*
by (*metis lborel-prod*)
qed

lemma *gauge-integral-Fubini-curve-bounded-region-x*:

fixes $f\ g :: 'a::euclidean-space * 'b::euclidean-space \Rightarrow 'c::euclidean-space$ **and**
 $g1\ g2 :: 'a \Rightarrow 'b$ **and**
 $s :: ('a * 'b)\ set$

assumes *fun-lesbegue-integrable: integrable lborel f* **and**
x-axis-gauge-integrable: $\bigwedge x.\ (\lambda y.\ f(x,\ y))\ integrable-on\ UNIV$ **and**

x-axis-integral-measurable: $(\lambda x.\ integral\ UNIV\ (\lambda y.\ f(x,\ y))) \in borel-measurable\ lborel$ **and**

f-is-g-indicator: $f = (\lambda x.\ if\ x \in s\ then\ g\ x\ else\ 0)$ **and**

s-is-bounded-by-g1-and-g2: $s = \{(x,y).\ (\forall i \in Basis.\ a \cdot i \leq x \cdot i \wedge x \cdot i \leq b \cdot i)$

\wedge

$(\forall i \in Basis.\ (g1\ x) \cdot i \leq y \cdot i \wedge y \cdot i \leq (g2\ x) \cdot i)$

shows *integral s g = integral (cbox a b) (\lambda x. integral (cbox (g1 x) (g2 x)) (\lambda y. g(x,y)))*

proof –

have *two-D-integral-to-one-D: integral UNIV f = integral UNIV (\lambda x. integral UNIV (\lambda y. f(x,y)))*

using *gauge-integral-Fubini-universe-x* **and** *fun-lesbegue-integrable* **and** *x-axis-integral-measurable*
by *auto*

have *one-d-integral-integrable: (\lambda x. integral UNIV (\lambda y. f(x,y))) integrable-on UNIV*

using *gauge-integral-Fubini-universe-x(2)* **and** *assms*

by *blast*

have *case-x-in-range*:

$\forall x \in cbox\ a\ b.\ integral\ (cbox\ (g1\ x)\ (g2\ x))\ (\lambda y.\ g(x,y)) = integral\ UNIV\ (\lambda y.\ f(x,y))$

proof

fix $x :: 'a$

assume *within-range: $x \in (cbox\ a\ b)$*

let *?f-one-D-spec = $(\lambda y.\ if\ y \in (cbox\ (g1\ x)\ (g2\ x))\ then\ g(x,y)\ else\ 0)$*

have *f-one-D-region: $(\lambda y.\ f(x,y)) = (\lambda y.\ if\ y \in cbox\ (g1\ x)\ (g2\ x)\ then\ g(x,y)\ else\ 0)$*

proof

fix $y :: 'b$

show $f\ (x,\ y) = (if\ y \in (cbox\ (g1\ x)\ (g2\ x))\ then\ g\ (x,\ y)\ else\ 0)$

using *within-range*

by (*force simp add: cbox-def f-is-g-indicator s-is-bounded-by-g1-and-g2*)

qed
have *zero-out-of-bound*: $\forall y. y \notin \text{cbox } (g1\ x) (g2\ x) \longrightarrow f(x,y) = 0$
using *f-is-g-indicator* **and** *s-is-bounded-by-g1-and-g2*
by (*auto simp add: cbox-def*)
have $(\lambda y. f(x,y))$ *integrable-on cbox* $(g1\ x) (g2\ x)$
proof –
have *?f-one-D-spec integrable-on UNIV*
using *f-one-D-region* **and** *x-axis-gauge-integrable*
by *metis*
then have *?f-one-D-spec integrable-on cbox* $(g1\ x) (g2\ x)$
using *integrable-on-subcbox* **by** *blast*
then show *?thesis* **using** *f-one-D-region* **by** *auto*
qed
then have *f-integrale-x*: $((\lambda y. f(x,y)) \text{ has-integral } (\text{integral } (\text{cbox } (g1\ x) (g2\ x)) (\lambda y. f(x,y)))) (\text{cbox } (g1\ x) (g2\ x))$
using *integrable-integral* **and** *within-range* **and** *x-axis-gauge-integrable*
by *auto*
have $\text{integral } (\text{cbox } (g1\ x) (g2\ x)) (\lambda y. f(x,y)) = \text{integral UNIV } (\lambda y. f(x,y))$
using *has-integral-on-superset*[*OF f-integrale-x - Set.subset-UNIV*] *zero-out-of-bound*
by (*simp add: integral-unique*)
then have $((\lambda y. f(x,y)) \text{ has-integral } \text{integral UNIV } (\lambda y. f(x,y))) (\text{cbox } (g1\ x) (g2\ x))$
using *f-integrale-x*
by *simp*
then have $((\lambda y. g(x,y)) \text{ has-integral } \text{integral UNIV } (\lambda y. f(x,y))) (\text{cbox } (g1\ x) (g2\ x))$
by (*simp add: f-one-D-region*)
then show $\text{integral } (\text{cbox } (g1\ x) (g2\ x)) (\lambda y. g(x,y)) = \text{integral UNIV } (\lambda y. f(x,y))$
by *auto*
qed
have *case-x-not-in-range*:
 $\forall x. x \notin \text{cbox } a\ b \longrightarrow \text{integral UNIV } (\lambda y. f(x,y)) = 0$
proof
fix $x::'a$
have $x \notin (\text{cbox } a\ b) \longrightarrow (\forall y. f(x,y) = 0)$
by (*auto simp add: s-is-bounded-by-g1-and-g2 f-is-g-indicator cbox-def*)
then show $x \notin \text{cbox } a\ b \longrightarrow \text{integral UNIV } (\lambda y. f(x,y)) = 0$
by (*simp*)
qed
have *RHS*: $\text{integral UNIV } (\lambda x. \text{integral UNIV } (\lambda y. f(x,y))) = \text{integral } (\text{cbox } a\ b) (\lambda x. \text{integral } (\text{cbox } (g1\ x) (g2\ x)) (\lambda y. g(x,y)))$
proof –
let *?first-integral* = $(\lambda x. \text{integral } (\text{cbox } (g1\ x) (g2\ x)) (\lambda y. g(x,y)))$
let *?x-integral-cases* = $(\lambda x. \text{if } x \in \text{cbox } a\ b \text{ then } ?\text{first-integral } x \text{ else } 0)$
have *x-integral-cases-integral*: $(\lambda x. \text{integral UNIV } (\lambda y. f(x,y))) = ?\text{x-integral-cases}$
using *case-x-in-range* **and** *case-x-not-in-range*
by *auto*

have ((λx . *integral UNIV* (λy . $f(x,y)$)) *has-integral* (*integral UNIV* f)) *UNIV*
using *two-D-integral-to-one-D one-d-integral-integrable* **by** *auto*
then have (*?x-integral-cases has-integral* (*integral UNIV* f)) *UNIV*
using *x-integral-cases-integral* **by** *auto*
then have (*?first-integral has-integral* (*integral UNIV* f)) (*cbox a b*)
using *has-integral-restrict-UNIV*[*of cbox a b ?first-integral integral UNIV f*]
by *auto*
then show *?thesis*
using *two-D-integral-to-one-D* **by** (*simp add: integral-unique*)
qed
have *f-integrable:f integrable-on UNIV*
using *fun-lesbegue-integrable and integrable-on-lborel*
by *auto*
then have *LHS: integral UNIV f = integral s g*
using *assms(4) integrable-integral* **by** *fastforce*
then show *?thesis*
using *RHS and two-D-integral-to-one-D*
by *auto*
qed

lemma *gauge-integral-Fubini-curve-bounded-region-y:*
fixes $f g :: ('a::euclidean-space * 'b::euclidean-space) \Rightarrow 'c::euclidean-space$ **and**
 $g1 g2 :: 'b \Rightarrow 'a$ **and**
 $s :: ('a * 'b) set$
assumes *fun-lesbegue-integrable: integrable lborel f* **and**
y-axis-gauge-integrable: $\bigwedge x. (\lambda y. f(y, x))$ integrable-on UNIV **and**

y-axis-integral-measurable: $(\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) \in \text{borel-measurable lborel}$ **and**
f-is-g-indicator: $f = (\lambda x. \text{if } x \in s \text{ then } g x \text{ else } 0)$ **and**
s-is-bounded-by-g1-and-g2: $s = \{(y, x). (\forall i \in \text{Basis}. a \cdot i \leq x \cdot i \wedge x \cdot i \leq b \cdot i) \wedge$
 $(\forall i \in \text{Basis}. (g1 x) \cdot i \leq y \cdot i \wedge y \cdot i \leq$
 $(g2 x) \cdot i)\}$
shows *integral s g = integral (cbox a b) ($\lambda x. \text{integral (cbox (g1 x) (g2 x)) } (\lambda y. g(y, x))$)*
proof –
have *two-D-integral-to-one-D: integral UNIV f = integral UNIV ($\lambda x. \text{integral UNIV } (\lambda y. f(y, x))$)*
using *gauge-integral-Fubini-universe-y and fun-lesbegue-integrable and y-axis-integral-measurable*
by *auto*
have *one-d-integral-integrable: ($\lambda x. \text{integral UNIV } (\lambda y. f(y, x))$) integrable-on UNIV*
using *gauge-integral-Fubini-universe-y(2) and assms*
by *blast*
have *case-y-in-range:*
 $\forall x \in \text{cbox } a \text{ b. integral (cbox (g1 x) (g2 x)) } (\lambda y. g(y, x)) = \text{integral UNIV } (\lambda y. f(y, x))$
proof

```

fix x::'b
assume within-range:  $x \in (\text{cbox } a \ b)$ 
let ?f-one-D-spec =  $(\lambda y. \text{if } y \in (\text{cbox } (g1 \ x) \ (g2 \ x)) \text{ then } g(y, x) \text{ else } 0)$ 
have f-one-D-region:  $(\lambda y. f(y, x)) = (\lambda y. \text{if } y \in \text{cbox } (g1 \ x) \ (g2 \ x) \text{ then } g(y,$ 
x) else 0)
proof
  fix y::'a
  show  $f(y, x) = (\text{if } y \in (\text{cbox } (g1 \ x) \ (g2 \ x)) \text{ then } g(y, x) \text{ else } 0)$ 
    using within-range
    by (force simp add: cbox-def f-is-g-indicator s-is-bounded-by-g1-and-g2)
qed
have zero-out-of-bound:  $\forall y. y \notin \text{cbox } (g1 \ x) \ (g2 \ x) \longrightarrow f(y, x) = 0$ 
  using f-is-g-indicator and s-is-bounded-by-g1-and-g2
  by (auto simp add: cbox-def)
have  $(\lambda y. f(y, x))$  integrable-on  $\text{cbox } (g1 \ x) \ (g2 \ x)$ 
proof -
  have ?f-one-D-spec integrable-on UNIV
    using f-one-D-region and y-axis-gauge-integrable
    by metis
  then have ?f-one-D-spec integrable-on  $\text{cbox}(g1 \ x) \ (g2 \ x)$ 
    using integrable-on-subcbox
    by blast
  then show ?thesis using f-one-D-region by auto
qed
then have f-integrale-y:  $((\lambda y. f(y, x)) \text{ has-integral } (\text{integral } (\text{cbox } (g1 \ x) \ (g2$ 
x))  $(\lambda y. f(y, x))))$   $(\text{cbox } (g1 \ x) \ (g2 \ x))$ 
  using integrable-integral and within-range and y-axis-gauge-integrable
  by auto
have  $\text{integral } (\text{cbox } (g1 \ x) \ (g2 \ x)) (\lambda y. f(y, x)) = \text{integral UNIV } (\lambda y. f(y,$ 
x))
  using has-integral-on-superset[OF f-integrale-y - Set.subset-UNIV] zero-out-of-bound
  by (simp add: integral-unique)
then have  $((\lambda y. f(y, x)) \text{ has-integral } \text{integral UNIV } (\lambda y. f(y, x)))$   $(\text{cbox } (g1$ 
x)  $(g2 \ x))$ 
  using f-integrale-y
  by simp
then have  $((\lambda y. g(y, x)) \text{ has-integral } \text{integral UNIV } (\lambda y. f(y, x)))$   $(\text{cbox } (g1$ 
x)  $(g2 \ x))$ 
  using f-one-D-region by fastforce
then show  $\text{integral } (\text{cbox } (g1 \ x) \ (g2 \ x)) (\lambda y. g(y, x)) = \text{integral UNIV } (\lambda y. f$ 
(y, x))
  by auto
qed
have case-y-not-in-range:
   $\forall x. x \notin \text{cbox } a \ b \longrightarrow \text{integral UNIV } (\lambda y. f(y, x)) = 0$ 
proof
  fix x::'b
  have  $x \notin (\text{cbox } a \ b) \longrightarrow (\forall y. f(y, x) = 0)$ 
    apply (simp add: s-is-bounded-by-g1-and-g2 f-is-g-indicator cbox-def)

```



```

    by auto
  then show  $x \notin \text{cbox } a \ b \longrightarrow \text{integral UNIV } (\lambda y. f(y, x)) = 0$ 
    by (simp)
qed
have RHS:  $\text{integral UNIV } (\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) = \text{integral } (\text{cbox } a \ b) (\lambda x. \text{integral } (\text{cbox } (g1 \ x) \ (g2 \ x)) (\lambda y. g(y, x)))$ 
proof -
  let ?first-integral =  $(\lambda x. \text{integral } (\text{cbox } (g1 \ x) \ (g2 \ x)) (\lambda y. g(y, x)))$ 
  let ?x-integral-cases =  $(\lambda x. \text{if } x \in \text{cbox } a \ b \ \text{then } ?\text{first-integral } x \ \text{else } 0)$ 
  have y-integral-cases-integral:  $(\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) = ?x\text{-integral-cases}$ 
    using case-y-in-range and case-y-not-in-range
    by auto
  have  $((\lambda x. \text{integral UNIV } (\lambda y. f(y, x))) \text{ has-integral } (\text{integral UNIV } f)) \text{ UNIV}$ 
    using two-D-integral-to-one-D
      one-d-integral-integrable
    by auto
  then have  $(?x\text{-integral-cases} \text{ has-integral } (\text{integral UNIV } f)) \text{ UNIV}$ 
    using y-integral-cases-integral by auto
  then have  $(?first\text{-integral} \text{ has-integral } (\text{integral UNIV } f)) (\text{cbox } a \ b)$ 
    using has-integral-restrict-UNIV[of cbox a b ?first-integral integral UNIV f]
    by auto
  then show ?thesis
    using two-D-integral-to-one-D
    by (simp add: integral-unique)
qed
have f-integrable:f integrable-on UNIV
  using fun-lesbegue-integrable and integrable-on-lborel
  by auto
then have LHS:  $\text{integral UNIV } f = \text{integral } s \ g$ 
  apply (simp add: f-is-g-indicator)
  using integrable-restrict-UNIV
    integral-restrict-UNIV
  by auto
then show ?thesis
  using RHS and two-D-integral-to-one-D
  by auto
qed

lemma gauge-integral-by-substitution:
  fixes f::(real  $\Rightarrow$  real) and
    g::(real  $\Rightarrow$  real) and
    g'::real  $\Rightarrow$  real and
    a::real and
    b::real
  assumes a-le-b:  $a \leq b$  and
    ga-le-gb:  $g \ a \leq g \ b$  and
    g'-derivative:  $\forall x \in \{a..b\}. (g \text{ has-vector-derivative } (g' \ x)) \text{ (at } x \ \text{within } \{a..b\})$ 
  and
    g'-continuous: continuous-on  $\{a..b\}$  g' and

```

f-continuous: continuous-on $(g \text{ ‘ } \{a..b\}) f$
shows $\text{integral } \{g \ a..g \ b\} (f) = \text{integral } \{a..b\} (\lambda x. f(g \ x) * (g' \ x))$
proof –
have $\forall x \in \{a..b\}. (g \text{ has-real-derivative } (g' \ x)) \text{ (at } x \text{ within } \{a..b\})$
using *has-real-derivative-iff-has-vector-derivative*[of *g*] **and** *g'-derivative*
by *auto*
then have 2: $\text{interval-lebesgue-integral lborel } (ereal \ (a)) \ (ereal \ (b)) \ (\lambda x. g' \ x *_{\mathbb{R}} f \ (g \ x))$
 $= \text{interval-lebesgue-integral lborel } (ereal \ (g \ a)) \ (ereal \ (g \ b)) \ f$
using *interval-integral-substitution-finite*[of *a b g g' f*] **and** *g'-continuous* **and**
a-le-b **and** *f-continuous*
by *auto*
have *g-continuous: continuous-on* $\{a .. b\} \ g$
using *Derivative.differentiable-imp-continuous-on*
apply (*simp add: differentiable-on-def differentiable-def*)
by (*metis continuous-on-vector-derivative g'-derivative*)
have *set-integrable lborel* $\{a .. b\} (\lambda x. g' \ x *_{\mathbb{R}} f \ (g \ x))$
proof –
have *continuous-on* $\{a .. b\} (\lambda x. g' \ x *_{\mathbb{R}} f \ (g \ x))$
proof –
have *continuous-on* $\{a .. b\} (\lambda x. f \ (g \ x))$
proof –
show *?thesis*
using *Topological-Spaces.continuous-on-compose f-continuous g-continuous*
by *auto*
qed
then show *?thesis*
using *Limits.continuous-on-mult g'-continuous*
by *auto*
qed
then show *?thesis*
using *borel-integrable-atLeastAtMost'* **by** *blast*
qed
then have 0: $\text{interval-lebesgue-integral lborel } (ereal \ (a)) \ (ereal \ (b)) \ (\lambda x. g' \ x *_{\mathbb{R}} f \ (g \ x))$
 $= \text{integral } \{a .. b\} (\lambda x. g' \ x *_{\mathbb{R}} f \ (g \ x))$
using *a-le-b* **and** *interval-integral-eq-integral*
by (*metis (no-types)*)
have *set-integrable lborel* $\{g \ a .. g \ b\} f$
proof –
have *continuous-on* $\{g \ a .. g \ b\} f$
proof –
have $\{g \ a .. g \ b\} \subseteq g \text{ ‘ } \{a .. b\}$
using *g-continuous*
by (*metis a-le-b atLeastAtMost-iff atLeastatMost-subset-iff continuous-image-closed-interval imageI order-refl*)
then show *continuous-on* $\{g \ a .. g \ b\} f$
using *f-continuous continuous-on-subset*
by *blast*

```

qed
then show ?thesis
  using borel-integrable-atLeastAtMost'
  by blast
qed
then have 1: interval-lebesgue-integral lborel (ereal (g a)) (ereal (g b)) f
  = integral {g a .. g b} f
  using ga-le-gb and interval-integral-eq-integral
  by (metis (no-types))
then show ?thesis
  using 0 and 1 and 2
  by (metis (no-types, lifting) Henstock-Kurzweil-Integration.integral-cong mult.commute
real-scaleR-def)
qed

```

```

lemma frontier-ic:
  assumes a < (b::real)
  shows frontier {a<..b} = {a,b}
  apply(simp add: frontier-def)
  using assms
  by auto

```

```

lemma frontier-ci:
  assumes a < (b::real)
  shows frontier {a<..} = {a,b}
  apply(simp add: frontier-def)
  using assms
  by auto

```

```

lemma ic-not-closed:
  assumes a < (b::real)
  shows ¬ closed {a<..b}
  using assms frontier-subset-eq frontier-ic greaterThanAtMost-iff by blast

```

```

lemma closure-ic-union-ci:
  assumes a < (b::real) b < c
  shows closure ({a..} ∪ {b<..

```

```

lemma interior-ic-ci-union:
  assumes a < (b::real) b < c
  shows b ∉ (interior ({a..} ∪ {b<..} ∪ {b<..

```

lemma *frontier-ic-union-ci*:
assumes $a < (b::real) < c$
shows $b \in \text{frontier} (\{a..<b\} \cup \{b<..c\})$
using *closure-ic-union-ci* *assms interior-ic-ci-union*
by (*simp add: frontier-def*)

lemma *ic-union-ci-not-closed*:
assumes $a < (b::real) < c$
shows $\neg \text{closed} (\{a..<b\} \cup \{b<..c\})$
proof –
have $b \notin (\{a..<b\} \cup \{b<..c\})$ **by** *auto*
then show *?thesis*
using *assms frontier-subset-eq frontier-ic-union-ci[OF assms]*
by (*auto simp only: subset-iff*)
qed

lemma *integrable-continuous-*:
fixes $f :: 'b::euclidean-space \Rightarrow 'a::banach$
assumes *continuous-on* (*cbox a b*) f
shows f *integrable-on* *cbox a b*
by (*simp add: assms integrable-continuous*)

lemma *removing-singletons-from-div*:
assumes $\forall t \in S. \exists c d :: real. c < d \wedge \{c..d\} = t$
 $\{x\} \cup \bigcup S = \{a..b\} \quad a < x < b$
finite S
shows $\exists t \in S. x \in t$
proof (*rule ccontr*)
assume $\neg (\exists t \in S. x \in t)$
then have $\forall t \in S. x \notin t$ **by** *auto*
then have $x \notin \bigcup S$ **by** *auto*
then have $i: \bigcup S = \{a..b\} - \{x\}$ **using** *assms (2)* **by** *auto*
have $x \in \{a..b\}$ **using** *assms* **by** *auto*
then have $\{a..b\} - \{x\} = \{a..<x\} \cup \{x<..b\}$ **by** *auto*
then have $0: \bigcup S = \{a..<x\} \cup \{x<..b\}$ **using** i **by** *auto*
have $1: \text{closed} (\bigcup S)$
apply (*rule closed-Union*)
proof –
show *finite S*
using *assms* **by** *auto*
show $\forall T \in S. \text{closed } T$ **using** *assms* **by** *auto*
qed
show *False* **using** $0\ 1$ *ic-union-ci-not-closed* *assms* **by** *auto*
qed

lemma *remove-singleton-from-division-of*:
assumes A *division-of* $\{a::real..b\} \quad a < b$
assumes $x \in \{a..b\}$

shows $\exists c d. c < d \wedge \{c..d\} \in A \wedge x \in \{c..d\}$
proof –
from *assms* **have** $x \text{ islimpt } \{a..b\}$
by (*intro connected-imp-perfect*) *auto*
also have $\{a..b\} = \{x. \{x..x\} \in A\} \cup (\{a..b\} - \{x. \{x..x\} \in A\})$
using *assms* **by** *auto*
also have $x \text{ islimpt } \dots \longleftrightarrow x \text{ islimpt } \{a..b\} - \{x. \{x..x\} \in A\}$
proof (*intro islimpt-Un-finite*)
have $\{x. \{x..x\} \in A\} \subseteq \text{Inf } 'A$
proof *safe*
fix x **assume** $\{x..x\} \in A$
thus $x \in \text{Inf } 'A$
by (*auto intro!: bexI[of - {x}] simp: image-iff*)
qed
moreover from *assms* **have** *finite A* **by** (*auto simp: division-of-def*)
hence *finite (Inf 'A)* **by** *auto*
ultimately show *finite {x. {x..x} \in A}* **by** (*rule finite-subset*)
qed
also have $\{a..b\} = \bigcup A$
using *assms* **by** (*auto simp: division-of-def*)
finally have $x \text{ islimpt } \bigcup (A - \text{range } (\lambda x. \{x..x\}))$
by (*rule islimpt-subset*) *auto*
moreover have *closed* $(\bigcup (A - \text{range } (\lambda x. \{x..x\})))$
using *assms* **by** (*intro closed-Union*) *auto*
ultimately have $x \in (\bigcup (A - \text{range } (\lambda x. \{x..x\})))$
by (*auto simp: closed-limpt*)
then obtain X **where** $x \in X \ X \in A \ X \notin \text{range } (\lambda x. \{x..x\})$
by *blast*
moreover from *division-ofD(2)[OF assms(1) this(2)] division-ofD(3)[OF assms(1) this(2)]*
division-ofD(4)[OF assms(1) this(2)]
obtain $c d$ **where** $X = \text{cbox } c d \ X \subseteq \{a..b\} \ X \neq \{\}$ **by** *blast*
ultimately have $c \leq d$ **by** *auto*
have $c \neq d$
proof
assume $c = d$
with $\langle X = \text{cbox } c d \rangle$ **have** $X = \{c..c\}$ **by** *auto*
hence $X \in \text{range } (\lambda x. \{x..x\})$ **by** *blast*
with $\langle X \notin \text{range } (\lambda x. \{x..x\}) \rangle$ **show** *False* **by** *contradiction*
qed
with $\langle c \leq d \rangle$ **have** $c < d$ **by** *simp*
with $\langle X = \text{cbox } c d \rangle$ **and** $\langle x \in X \rangle$ **and** $\langle X \in A \rangle$ **show** *?thesis*
by *auto*
qed
lemma *remove-singleton-from-tagged-division-of:*
assumes *A tagged-division-of {a::real..b} a < b*
assumes $x \in \{a..b\}$
shows $\exists k c d. c < d \wedge (k, \{c..d\}) \in A \wedge x \in \{c..d\}$

using *remove-singleton-from-division-of*[*OF division-of-tagged-division*[*OF assms(1)*]
assms(2)]

using *assms(3)* **by** *fastforce*

lemma *tagged-div-wo-singletons*:

assumes p *tagged-division-of* $\{a::\text{real}..b\}$ $a < b$

shows $(p - \{xk. \exists x y. xk = (x, \{y\})\})$ *tagged-division-of* *cbox* a b

using *remove-singleton-from-tagged-division-of*[*OF assms*] *assms*

apply(*auto simp add: tagged-division-of-def tagged-partial-division-of-def*)

apply *blast*

apply *blast*

apply *blast*

by *fastforce*

lemma *tagged-div-wo-empty*:

assumes p *tagged-division-of* $\{a::\text{real}..b\}$ $a < b$

shows $(p - \{xk. \exists x. xk = (x, \{\})\})$ *tagged-division-of* *cbox* a b

using *remove-singleton-from-tagged-division-of*[*OF assms*] *assms*

apply(*auto simp add: tagged-division-of-def tagged-partial-division-of-def*)

apply *blast*

apply *blast*

apply *blast*

by *fastforce*

lemma *fine-diff*:

assumes γ *fine* p

shows γ *fine* $(p - s)$

apply (*auto simp add: fine-def*)

using *assms* **by** *auto*

lemma *tagged-div-tage-notin-set*:

assumes *finite* $(s::\text{real set})$

p *tagged-division-of* $\{a..b\}$

γ *fine* p $(\forall (x, K) \in p. \exists c d::\text{real}. c < d \wedge K = \{c..d\})$ *gauge* γ

shows $\exists p' \gamma'. p'$ *tagged-division-of* $\{a..b\} \wedge$

γ' *fine* $p' \wedge (\forall (x, K) \in p'. x \notin s) \wedge$ *gauge* γ'

proof –

have $(\forall (x::\text{real}, K) \in p. \exists x'. x' \notin s \wedge x' \in \text{interior } K)$

proof –

{fix $x::\text{real}$

fix K

assume $\text{ass}: (x::\text{real}, K) \in p$

have $(\forall (x, K) \in p. \text{infinite } (\text{interior } K))$

using *assms(4)* *infinite-Ioo interior-atLeastAtMost-real*

by (*smt (verit) split-beta*)

then have $i: \text{infinite } (\text{interior } K)$ **using** *ass* **by** *auto*

have $\exists x'. x' \notin s \wedge x' \in \text{interior } K$

using *infinite-imp-nonempty*[*OF Diff-infinite-finite*[*OF assms(1) i*]] **by** *auto*}

then show *?thesis* **by** *auto*

```

qed
then obtain  $f$  where  $f: (\forall (x::real, K)\in p. (f (x,K)) \notin s \wedge (f (x,K)) \in interior$ 
 $K)$ 
  using choice-iff [where  $?Q = \lambda(x,K) x'. (x::real, K)\in p \longrightarrow x' \notin s \wedge x' \in$ 
interior K]
  apply (auto simp add: case-prod-beta)
  by metis
have  $f': (\forall (x::real, K)\in p. (f (x,K)) \notin s \wedge (f (x,K)) \in K)$ 
  using  $f$  interior-subset
  by (auto simp add: case-prod-beta subset-iff)
let  $?p' = \{m. (\exists xK. m = ((f xK), snd xK) \wedge xK \in p)\}$ 
have  $0: (\forall (x, K)\in ?p'. x \notin s)$ 
  using  $f$ 
  by (auto simp add: case-prod-beta)
have  $i: finite \{(f (a, b), b) \mid a b. (a, b) \in p\}$ 
proof -
  have  $a: \{(f (a, b), b) \mid a b. (a, b) \in p\} = (\% (a,b). (f(a,b),b)) ' p$  by auto
  have  $b: finite p$  using assms(2) by auto
  show  $?thesis$  using  $a b$  by auto
qed
have  $1: ?p'$  tagged-division-of  $\{a..b\}$ 
  using assms(2)  $f'$ 
apply (auto simp add: tagged-division-of-def tagged-partial-division-of-def case-prod-beta)
  apply (metis i)
  apply blast
  apply blast
by fastforce

have  $f\text{-inj}: inj\text{-on } f p$ 
  unfolding inj-on-def
proof (intro strip)
  fix  $x y$ 
  assume  $x \in p y \in p f x = f y$ 
  then show  $x = y$ 
  using  $f$  tagged-division-ofD(5)[OF assms(2)]
  by (smt (verit, del-insts) IntI case-prodE empty-iff)
qed
let  $? \gamma' = \lambda x. if (\exists xK \in p. f xK = x) then (\gamma \circ fst \circ the\text{-inv-into } p f) x$  else  $\gamma$ 
 $x$ 
have  $2: ? \gamma'$  fine  $?p'$  using assms(3)
  by (force simp add: fine-def case-prod-beta the-inv-into-f-f[OF f-inj])
have  $3: gauge ? \gamma'$ 
  using assms(5) assms(3) f'
  by (force simp add: fine-def gauge-def case-prod-beta the-inv-into-f-f[OF f-inj])
have  $?p'$  tagged-division-of  $\{a..b\} \wedge ? \gamma'$  fine  $?p' \wedge (\forall (x, K)\in ?p'. x \notin s) \wedge gauge$ 
 $? \gamma'$ 
  using  $0 1 2 3$  by auto
then show  $?thesis$  by meson
qed

```

lemma *has-integral-bound-spike-finite*:
fixes $f :: 'a::\text{euclidean-space} \Rightarrow 'b::\text{real-normed-vector}$
assumes $0 \leq B$ **and** *finite S*
and $f: (f \text{ has-integral } i) (\text{cbox } a \text{ } b)$
and $leB: \bigwedge x. x \in \text{cbox } a \text{ } b - S \implies \text{norm } (f \ x) \leq B$
shows $\text{norm } i \leq B * \text{content } (\text{cbox } a \text{ } b)$
proof –
define g **where** $g \equiv (\lambda x. \text{if } x \in S \text{ then } 0 \text{ else } f \ x)$
then have $\bigwedge x. x \in \text{cbox } a \text{ } b - S \implies \text{norm } (g \ x) \leq B$
using leB **by** *simp*
moreover have $(g \text{ has-integral } i) (\text{cbox } a \text{ } b)$
using *has-integral-spike-finite* [*OF* $\langle \text{finite } S \rangle$ - f]
by (*simp add: g-def*)
ultimately show *?thesis*
by (*simp add:* $\langle 0 \leq B \rangle$ *g-def has-integral-bound*)

qed

lemma *has-integral-bound-*:
fixes $f :: \text{real} \Rightarrow 'a::\text{real-normed-vector}$
assumes $a < b$
and $0 \leq B$
and $f: (f \text{ has-integral } i) (\text{cbox } a \text{ } b)$
and *finite s*
and $\forall x \in (\text{cbox } a \text{ } b) - s. \text{norm } (f \ x) \leq B$
shows $\text{norm } i \leq B * \text{content } (\text{cbox } a \text{ } b)$
using *has-integral-bound-spike-finite* *assms* **by** *blast*

corollary *has-integral-bound-real'*:
fixes $f :: \text{real} \Rightarrow 'b::\text{real-normed-vector}$
assumes $0 \leq B$
and $f: (f \text{ has-integral } i) (\text{cbox } a \text{ } b)$
and *finite s*
and $\forall x \in (\text{cbox } a \text{ } b) - s. \text{norm } (f \ x) \leq B$
shows $\text{norm } i \leq B * \text{content } \{a..b\}$
by (*metis* *assms*(1) *assms*(3) *assms*(4) *box-real*(2) *f has-integral-bound-spike-finite*)

lemma *integral-has-vector-derivative-continuous-at'*:
fixes $f :: \text{real} \Rightarrow 'a::\text{banach}$
assumes *finite s*
and $f: f \text{ integrable-on } \{a..b\}$
and $x \in \{a..b\} - s$
and $fx: \text{continuous } (\text{at } x \text{ within } (\{a..b\} - s)) \ f$
shows $((\lambda u. \text{integral } \{a..u\} \ f) \text{ has-vector-derivative } f \ x) (\text{at } x \text{ within } (\{a..b\} - s))$
proof –
let $?I = \lambda a \ b. \text{integral } \{a..b\} \ f$
{ **fix** $e::\text{real}$
assume $e > 0$


```

obtain  $d$  where  $d > 0$  and  $d: \bigwedge x'. \llbracket x' \in \{a..b\} - s; |x' - x| < d \rrbracket \implies \text{norm}(f$ 
 $x' - f x) \leq e$ 
using  $\langle e > 0 \rangle$   $fx$  by (auto simp: continuous-within-eps-delta dist-norm less-imp-le)
have  $\text{norm} (\text{integral} \{a..y\} f - \text{integral} \{a..x\} f - (y-x) *_R f x) \leq e * |y -$ 
 $x|$ 
if  $y: y \in \{a..b\} - s$  and  $yx: |y - x| < d$  for  $y$ 
proof (cases y < x)
case False
have  $f$  integrable-on  $\{a..y\}$ 
using  $f y$  by (simp add: integrable-subinterval-real)
then have Idiff:  $?I a y - ?I a x = ?I x y$ 
using False x by (simp add: algebra-simps Henstock-Kurzweil-Integration.integral-combine)
have  $fx\text{-int}$ :  $((\lambda u. f u - f x) \text{ has-integral integral } \{x..y\} f - (y-x) *_R f x)$ 
 $\{x..y\}$ 
apply (rule has-integral-diff)
using  $x y$  apply (auto intro: integrable-integral [OF integrable-subinterval-real
 $[OF f]]$ )
using has-integral-const-real [of f x x y] False
apply simp
done
show ?thesis
using False
apply (simp add: abs-eq-content del: content-real-if measure-lborel-Icc)
apply (rule has-integral-bound-real'[where f=(\lambda u. f u - f x)])
using  $yx$  False d x y  $\langle e > 0 \rangle$  apply (auto simp add: Idiff fx-int)
proof-
let  $?M48 = \text{mset-set } s$ 
show  $\bigwedge z. y - x < d \implies (\bigwedge x'. a \leq x' \wedge x' \leq b \wedge x' \notin s \implies |x' - x| < d$ 
 $\implies \text{norm} (f x' - f x) \leq e) \implies 0 < e \implies z \notin \# ?M48 \implies a \leq x \implies x \notin s \implies$ 
 $y \leq b \implies y \notin s \implies x \leq z \implies z \leq y \implies \text{norm} (f z - f x) \leq e$ 
using assms by auto
qed
next
case True
have  $f$  integrable-on  $\{a..x\}$ 
using  $f x$  by (simp add: integrable-subinterval-real)
then have Idiff:  $?I a x - ?I a y = ?I y x$ 
using True x y by (simp add: algebra-simps Henstock-Kurzweil-Integration.integral-combine)
have  $fx\text{-int}$ :  $((\lambda u. f u - f x) \text{ has-integral integral } \{y..x\} f - (x - y) *_R f x)$ 
 $\{y..x\}$ 
apply (rule has-integral-diff)
using  $x y$  apply (auto intro: integrable-integral [OF integrable-subinterval-real
 $[OF f]]$ )
using has-integral-const-real [of f x y x] True
by simp
have  $\text{norm} (\text{integral} \{a..x\} f - \text{integral} \{a..y\} f - (x - y) *_R f x) \leq e * |y$ 
 $- x|$ 
using True
apply (simp add: abs-eq-content del: content-real-if measure-lborel-Icc)

```

```

apply (rule has-integral-bound-real'[where f=( $\lambda u. f u - f x$ )])
using yx True d x y <e>0> apply (auto simp add: Idiff fux-int)
proof -
  let ?M44 = mset-set s
  show  $\bigwedge xa. x - y < d \implies y < x \implies (\bigwedge x'. a \leq x' \wedge x' \leq b \wedge x' \notin s \implies$ 
 $|x' - x| < d \implies \text{norm } (f x' - f x) \leq e) \implies 0 < e \implies xa \notin \# ?M44 \implies x \leq b$ 
 $\implies x \notin s \implies a \leq y \implies y \notin s \implies y \leq xa \implies xa \leq x \implies \text{norm } (f xa - f x) \leq e$ 
  using assms by auto
  qed
  then show ?thesis
  by (simp add: algebra-simps norm-minus-commute)
  qed
  then have  $\exists d > 0. \forall y \in \{a..b\} - s. |y - x| < d \implies \text{norm } (\text{integral } \{a..y\} f -$ 
 $\text{integral } \{a..x\} f - (y-x) *_R f x) \leq e * |y - x|$ 
  using <d>0> by blast
}
then show ?thesis
by (simp add: has-vector-derivative-def has-derivative-within-alt bounded-linear-scaleR-left)
qed

```

lemma at-within-closed-interval-finite:

```

fixes x::real
assumes a < x x < b x  $\notin$  S finite S
shows (at x within {a..b} - S) = at x
proof -
  have interior ({a..b} - S) = {a<..<b} - S
  using <finite S>
  by (simp add: interior-diff finite-imp-closed)
  then show ?thesis
  using at-within-interior assms by fastforce
qed

```

lemma fundamental-theorem-of-calculus-interior-stronger':

```

fixes f :: real  $\Rightarrow$  'a::banach
assumes finite S
  and a  $\leq$  b  $\bigwedge x. x \in \{a <..<b\} - S \implies (f \text{ has-vector-derivative } f'(x))$  (at x
  within {a..b} - S)
  and continuous-on {a .. b} f
shows (f' has-integral (f b - f a)) {a .. b}
using assms fundamental-theorem-of-calculus-interior-strong at-within-cbox-finite
by (metis DiffD1 DiffD2 interior-atLeastAtMost-real interior-cbox interval-cbox)

```

lemma has-integral-substitution-general:

```

fixes f :: real  $\Rightarrow$  'a::euclidean-space and g :: real  $\Rightarrow$  real
assumes s: finite s and le: a  $\leq$  b
  and subset: g ' {a..b}  $\subseteq$  {c..d}
  and f: f integrable-on {c..d} continuous-on ({c..d} - (g ' s)) f
  and g : continuous-on {a..b} g inj-on g ({a..b}  $\cup$  s)
  and deriv [derivative-intros]:

```

$\bigwedge x. x \in \{a..b\} - s \implies (g \text{ has-field-derivative } g' x) \text{ (at } x \text{ within } \{a..b\})$
shows $((\lambda x. g' x *_R f (g x)) \text{ has-integral } (\text{integral } \{g a..g b\} f - \text{integral } \{g b..g a\} f)) \{a..b\}$
proof –
let $?F = \lambda x. \text{integral } \{c..g x\} f$
have *cont-int: continuous-on* $\{a..b\} ?F$
by $(\text{rule continuous-on-compose2}[OF - g(1) \text{ subset}] \text{ indefinite-integral-continuous-1 } f)+$
have *deriv*: $\bigwedge x. x \in \{a..b\} - s \implies (((\lambda x. \text{integral } \{c..x\} f) \circ g) \text{ has-vector-derivative } g' x *_R f (g x))$
 $\text{ (at } x \text{ within } (\{a..b\} - s))$
apply $(\text{rule has-vector-derivative-eq-rhs})$
apply $(\text{rule vector-diff-chain-within})$
apply $(\text{subst has-real-derivative-iff-has-vector-derivative } [\text{symmetric}])$
proof –
fix $x::\text{real}$
assume *ass*: $x \in \{a..b\} - s$
let $?f'g = g' x$
have *i*: $\{a..b\} - s \subseteq \{a..b\}$ **by** *auto*
have *ii*: $(g \text{ has-vector-derivative } g' x) \text{ (at } x \text{ within } \{a..b\})$ **using** $\text{deriv}[OF \text{ ass}]$
by $(\text{simp only: has-real-derivative-iff-has-vector-derivative})$
show $(g \text{ has-real-derivative } ?f'g) \text{ (at } x \text{ within } \{a..b\} - s)$
using $\text{has-vector-derivative-within-subset}[OF \text{ ii } i]$
by $(\text{simp only: has-real-derivative-iff-has-vector-derivative})$
next
let $?g'g = f \circ g$
show $\bigwedge x. x \in \{a..b\} - s \implies ((\lambda x. \text{integral } \{c..x\} f) \text{ has-vector-derivative } ?g'g x)$
 $\text{ (at } (g x) \text{ within } g' (\{a..b\} - s))$
proof –
fix $x::\text{real}$
assume *ass*: $x \in \{a..b\} - s$
have *finite* $(g' s)$ **using** *s by auto*
then have *i*: $((\lambda x. \text{integral } \{c..x\} f) \text{ has-vector-derivative } f(g x)) \text{ (at } (g x) \text{ within } (\{c..d\} - g' s))$
proof $(\text{rule integral-has-vector-derivative-continuous-at'})$
show $f \text{ integrable-on } \{c..d\}$ **using** *f by auto*
show $g x \in \{c..d\} - g' s$ **using** *ass subset*
by $(\text{smt (verit) Diff-iff } g(2) \text{ inf-sup-ord}(4) \text{ inj-on-image-mem-iff subsetD sup-ge1})$
show *continuous* $(\text{at } (g x) \text{ within } \{c..d\} - g' s) f$
using $\langle g x \in \{c..d\} - g' s \rangle \text{ continuous-on-eq-continuous-within } f(2)$ **by**
blast
qed
have *ii*: $g' (\{a..b\} - s) \subseteq (\{c..d\} - g' s)$
using $\text{subset } g(2)$
by $(\text{simp add: image-subset-iff inj-on-image-mem-iff})$
then show $((\lambda x. \text{integral } \{c..x\} f) \text{ has-vector-derivative } ?g'g x) \text{ (at } (g x) \text{ within } g' (\{a..b\} - s))$
using $\text{has-vector-derivative-within-subset } i$ **by** *fastforce*

```

qed
show  $\bigwedge x. x \in \{a..b\} - s \implies g' x *_R ?g' \exists x = g' x *_R f (g x)$  by auto
qed
have deriv: (?F has-vector-derivative  $g' x *_R f (g x)$ )
  (at  $x$  within  $\{a..b\} - s$ ) if  $x \in \{a < .. < b\} - (s)$  for  $x$ 
  using deriv[of  $x$ ] that by (simp add: at-within-Icc-at o-def)
have (( $\lambda x. g' x *_R f (g x)$ ) has-integral (?F  $b - ?F a$ )  $\{a..b\}$ )
  using cont-int
  using fundamental-theorem-of-calculus-interior-stronger'[OF  $s$  le deriv]
  by blast
also
from subset have  $g x \in \{c..d\}$  if  $x \in \{a..b\}$  for  $x$  using that by blast
from this[of  $a$ ] this[of  $b$ ] le have  $cd: c \leq g a \ g b \leq d \ c \leq g b \ g a \leq d$  by auto
have integral  $\{c..g b\} f - integral \{c..g a\} f = integral \{g a..g b\} f - integral$ 
 $\{g b..g a\} f$ 
proof cases
  assume  $g a \leq g b$ 
  note  $le = le$  this
  from  $cd$  have  $integral \{c..g a\} f + integral \{g a..g b\} f = integral \{c..g b\} f$ 
  by (meson Henstock-Kurzweil-Integration.integral-combine atLeastatMost-subset-iff
 $f(1)$  integrable-on-subinterval  $le(2)$  order-refl)
  with  $le$  show ?thesis
  by (cases  $g a = g b$ ) (simp-all add: algebra-simps)
next
  assume less:  $\neg g a \leq g b$ 
  then have  $le: g a \geq g b$  by simp
  from  $cd$  have  $integral \{c..g b\} f + integral \{g b..g a\} f = integral \{c..g a\} f$ 
  by (meson Henstock-Kurzweil-Integration.integral-combine atLeastatMost-subset-iff
 $f(1)$  integrable-on-subinterval  $le$  order-refl)
  with less show ?thesis
  by (simp-all add: algebra-simps)
qed
finally show ?thesis .
qed

```

lemma *has-integral-substitution-general--:*

```

fixes  $f :: real \Rightarrow 'a::euclidean-space$  and  $g :: real \Rightarrow real$ 
assumes  $s$ : finite  $s$  and  $le: a \leq b$  and  $s$ -subset:  $s \subseteq \{a..b\}$ 
  and subset:  $g ' \{a..b\} \subseteq \{c..d\}$ 
  and  $f$ :  $f$  integrable-on  $\{c..d\}$  continuous-on  $(\{c..d\} - (g ' s)) f$ 
  and  $g$ : continuous-on  $\{a..b\}$   $g$  inj-on  $g \{a..b\}$ 
  and deriv [derivative-intros]:
 $\bigwedge x. x \in \{a..b\} - s \implies (g$  has-field-derivative  $g' x)$  (at  $x$  within  $\{a..b\}$ )
shows (( $\lambda x. g' x *_R f (g x)$ ) has-integral (integral  $\{g a..g b\} f - integral \{g b..g$ 
 $a\} f$ ))  $\{a..b\}$ 
  using  $s$ -subset has-integral-substitution-general-[OF  $s$  le subset  $f g(1) - deriv$ ]
  by (simp add:  $g(2)$  sup-absorb1)

```

lemma *has-integral-substitution-general-':*

```

fixes f :: real ⇒ 'a::euclidean-space and g :: real ⇒ real
assumes s: finite s and le: a ≤ b and s': finite s'
  and subset: g ` {a..b} ⊆ {c..d}
  and f: f integrable-on {c..d} continuous-on ({c..d} - s') f
  and g : continuous-on {a..b} g ∀ x ∈ s'. finite (g - ` {x}) surj-on s' g inj-on g
  ({a..b} ∪ ((s ∪ g - ` s')))
  and deriv [derivative-intros]:
    ∧ x. x ∈ {a..b} - s ⇒ (g has-field-derivative g' x) (at x within {a..b})
  shows ((λ x. g' x *R f (g x)) has-integral (integral {g a..g b} f - integral {g b..g
a} f)) {a..b}
proof -
  have a: g - ` s' = ∪ {t. ∃ x. t = g - ` {x} ∧ x ∈ s'}
  using s s' by blast
  have finite (∪ {t. ∃ x. t = g - ` {x} ∧ x ∈ s'}) using s'
  by (metis (no-types, lifting) ‹g - ` s' = ∪ {g - ` {x} | x. x ∈ s'}› finite-UN-I
g(2) vimage-eq-UN)
  then have 0: finite (s ∪ (g - ` s'))
  using a s by simp
  have 1: continuous-on ({c..d} - g ` (s ∪ g - ` s')) f
  using f(2) surj-on-image-vimage-eq
  by (metis Diff-mono Un-upper2 continuous-on-subset equalityE g(3) image-Un)
  have 2: (∧ x. x ∈ {a..b} - (s ∪ g - ` s') ⇒ (g has-real-derivative g' x) (at x
within {a..b}))
  using deriv by auto
  show ?thesis using has-integral-substitution-general-[OF 0 assms(2) subset f(1)
1 g(1) g(4) 2]
  by auto
qed

end
theory Paths
  imports Derivs General-Utills Integrals
begin

```

```

lemma reverse-subpaths-join:
  shows subpath 1 (1 / 2) p +++ subpath (1 / 2) 0 p = reversepath p
  using reversepath-subpath join-subpaths-middle pathfinish-subpath pathstart-subpath
reversepath-joinpaths
  by (metis (no-types, lifting))

```

```

definition line-integral:: ('a::euclidean-space ⇒ 'a::euclidean-space) ⇒ (('a) set)
⇒ (real ⇒ 'a) ⇒ real where
line-integral F basis g ≡ integral {0 .. 1} (λ x. ∑ b ∈ basis. (F(g x) · b) * (vector-derivative
g (at x within {0..1}) · b))

```

```

definition line-integral-exists where

```

line-integral-exists F basis $\gamma \equiv (\lambda x. \sum b \in \text{basis}. F(\gamma x) \cdot b * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot b)) \text{ integrable-on } \{0..1\}$

lemma *line-integral-on-pair-straight-path:*

fixes $F :: ('a :: \text{euclidean-space}) \Rightarrow 'a$ **and** $g :: \text{real} \Rightarrow \text{real}$ **and** γ

assumes *gamma-const:* $\forall x. \gamma(x) \cdot i = a$

and *gamma-smooth:* $\forall x \in \{0 .. 1\}. \gamma$ *differentiable at* x

shows $(\text{line-integral } F \{i\} \gamma) = 0$ *(line-integral-exists F {i} \gamma)*

proof *(simp add: line-integral-def)*

have $*$: $F(\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i) = 0$

if $0 \leq x \wedge x \leq 1$ **for** x

proof $-$

have $((\lambda x. \gamma(x) \cdot i) \text{ has-vector-derivative } 0) \text{ (at } x)$

using *vector-derivative-const-at[of a x]* **and** *gamma-const*

by *auto*

then have $(\text{vector-derivative } \gamma \text{ (at } x) \cdot i) = 0$

using *derivative-component-fun-component[of \gamma x i]*

and *gamma-smooth and that*

by *(simp add: vector-derivative-at)*

then have $(\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0 .. 1\}) \cdot i) = 0$

using *has-vector-derivative-at-within vector-derivative-at-within-ivl that*

by *(metis atLeastAtMost-iff gamma-smooth vector-derivative-works zero-less-one)*

then show *?thesis*

by *auto*

qed

then have $((\lambda x. F(\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i)) \text{ has-integral } 0) \{0..1\}$

using *has-integral-is-0[of {0 .. 1} (\lambda x. F(\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i))]*

by *auto*

then have $((\lambda x. F(\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i)) \text{ integrable-on } \{0..1\})$

by *auto*

then show *line-integral-exists F {i} \gamma* **by** *(auto simp add: line-integral-exists-def)*

show $\text{integral } \{0..1\} (\lambda x. F(\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i)) = 0$

using $*$ *has-integral-is-0[of {0 .. 1} (\lambda x. F(\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i))]*

by *auto*

qed

lemma *line-integral-on-pair-path-strong:*

fixes $F :: ('a :: \text{euclidean-space}) \Rightarrow ('a)$ **and**

$g :: \text{real} \Rightarrow 'a$ **and**

$\gamma :: (\text{real} \Rightarrow 'a)$ **and**

$i :: 'a$

assumes *i-norm-1: norm i = 1* **and**

g-orthogonal-to-i: \forall x. g(x) \cdot i = 0 **and**

*gamma-is-in-terms-of-i: \gamma = (\lambda x. f(x) *_R i + g(f(x)))* **and**

gamma-smooth: γ *piecewise-C1-differentiable-on* $\{0 .. 1\}$ **and**
g-continuous-on-f: *continuous-on* $(f \text{ ' } \{0..1\})$ g **and**
path-start-le-path-end: $(\text{pathstart } \gamma) \cdot i \leq (\text{pathfinish } \gamma) \cdot i$ **and**
field-i-comp-cont: *continuous-on* $(\text{path-image } \gamma)$ $(\lambda x. F x \cdot i)$
shows *line-integral* $F \{i\} \gamma$
 $= \text{integral } (\text{cbox } ((\text{pathstart } \gamma) \cdot i) ((\text{pathfinish } \gamma) \cdot i)) (\lambda f\text{-var. } (F (f\text{-var}$
 $*_R i + g(f\text{-var})) \cdot i))$
line-integral-exists $F \{i\} \gamma$
proof (*simp add: line-integral-def*)
obtain s **where** *gamma-differentiable: finite* $s (\forall x \in \{0 .. 1\} - s. \gamma$ *differentiable*
at $x)$
using *gamma-smooth*
by (*auto simp add: C1-differentiable-on-eq piecewise-C1-differentiable-on-def*)
then have *gamma-i-component-smooth*: $\forall x \in \{0 .. 1\} - s. (\lambda x. \gamma x \cdot i)$ *differentiable*
at x
by *auto*
have *field-cont-on-path*: *continuous-on* $((\lambda x. \gamma x \cdot i) \text{ ' } \{0..1\}) (\lambda f\text{-var. } F (f\text{-var}$
 $*_R i + g f\text{-var}) \cdot i)$
proof -
have $0: (\lambda x. \gamma x \cdot i) = f$
proof
fix x
show $\gamma x \cdot i = f x$
using *g-orthogonal-to-i i-norm-1*
by (*simp only: gamma-is-in-terms-of-i real-inner-class.inner-add-left g-orthogonal-to-i*
inner-scaleR-left inner-same-Basis norm-eq-1)
qed
show *?thesis*
unfolding 0
apply (*rule continuous-on-compose2* [*of* - $(\lambda x. F(x) \cdot i)$ $f \text{ ' } \{0..1\}$ $(\lambda x. x$
 $*_R i + g x)$]
field-i-comp-cont g-continuous-on-f field-i-comp-cont continuous-intros)+
by (*auto simp add: gamma-is-in-terms-of-i path-image-def*)
qed
have *path-start-le-path-end'*: $\gamma 0 \cdot i \leq \gamma 1 \cdot i$ **using** *path-start-le-path-end* **by**
(auto simp add: pathstart-def pathfinish-def)
have *gamm-cont*: *continuous-on* $\{0..1\}$ $(\lambda a. \gamma a \cdot i)$
apply(*rule continuous-on-inner*)
using *gamma-smooth*
apply (*simp add: piecewise-C1-differentiable-on-def*)
using *continuous-on-const* **by** *auto*
then obtain $c d$ **where** *cd*: $c \leq d (\lambda a. \gamma a \cdot i) \text{ ' } \{0..1\} = \{c..d\}$
by (*meson continuous-image-closed-interval zero-le-one*)
then have *subset-cd*: $(\lambda a. \gamma a \cdot i) \text{ ' } \{0..1\} \subseteq \{c..d\}$ **by** *auto*
have *field-cont-on-path-cd*:
continuous-on $\{c..d\}$ $(\lambda f\text{-var. } F (f\text{-var } *_R i + g f\text{-var}) \cdot i)$
using *field-cont-on-path cd* **by** *auto*
have *path-vector-deriv-line-integrals*:
 $\forall x \in \{0..1\} - s. ((\lambda x. \gamma x \cdot i)$ *has-vector-derivative*

$vector\text{-}derivative (\lambda x. \gamma x \cdot i) (at x)$
($at x$)

using $gamma\text{-}i\text{-component}\text{-smooth}$ **and** $derivative\text{-component}\text{-fun}\text{-component}$
and

$vector\text{-}derivative\text{-works}$
by $blast$

then have $\forall x \in \{0..1\} - s. ((\lambda x. \gamma x \cdot i) \text{ has}\text{-}vector\text{-}derivative$
 $vector\text{-}derivative (\lambda x. \gamma x \cdot i) (at x \text{ within}$
 $(\{0..1\})))$
 $(at x \text{ within } (\{0..1\}))$

using $has\text{-}vector\text{-}derivative\text{-at}\text{-within}$ $vector\text{-}derivative\text{-at}\text{-within}\text{-ivl}$
by $fastforce$

then have $has\text{-}int:((\lambda x. vector\text{-}derivative (\lambda x. \gamma x \cdot i) (at x \text{ within } \{0..1\}) *_R$
 $(F ((\gamma x \cdot i) *_R i + g (\gamma x \cdot i)) \cdot i)) \text{ has}\text{-}integral$
 $integral \{\gamma 0 \cdot i.. \gamma 1 \cdot i\} (\lambda f\text{-}var. F (f\text{-}var *_R i + g f\text{-}var) \cdot i)) \{0..1\}$

using $has\text{-}integral\text{-substitution}\text{-strong}[OF \text{ gamma}\text{-}differentiable(1) \text{ rel}\text{-}simps(44)$
 $path\text{-}start\text{-}le\text{-}path\text{-}end' \text{ subset}\text{-}cd \text{ field}\text{-}cont\text{-}on\text{-}path\text{-}cd \text{ gamm}\text{-}cont,$
 $of (\lambda x. vector\text{-}derivative (\lambda x. \gamma(x) \cdot i) (at x \text{ within } (\{0..1\})))$
 $gamma\text{-}is\text{-}in\text{-}terms\text{-}of\text{-}i$

by ($auto \text{ simp only: has}\text{-}real\text{-}derivative\text{-}iff\text{-}has\text{-}vector\text{-}derivative$)

then have $has\text{-}int':((\lambda x. (F(\gamma(x)) \cdot i) * (vector\text{-}derivative (\lambda x. \gamma(x) \cdot i) (at x$
 $\text{ within } (\{0..1\})))) \text{ has}\text{-}integral$
 $integral \{((pathstart \gamma) \cdot i)..((pathfinish \gamma) \cdot i)\} (\lambda f\text{-}var. F (f\text{-}var *_R i +$
 $g f\text{-}var) \cdot i)) \{0..1\}$

using $gamma\text{-}is\text{-}in\text{-}terms\text{-}of\text{-}i \text{ i}\text{-}norm\text{-}1$

apply ($auto \text{ simp add: pathstart}\text{-}def \text{ pathfinish}\text{-}def$)

apply ($simp \text{ only: real}\text{-}inner\text{-}class.\text{inner}\text{-}add\text{-}left \text{ inner}\text{-}not\text{-}same\text{-}Basis \text{ g}\text{-}orthogonal\text{-}to\text{-}i$
 $\text{ inner}\text{-}scaleR\text{-}left \text{ norm}\text{-}eq\text{-}1$)

by ($auto \text{ simp add: algebra}\text{-}simps$)

have $substitute:$
 $integral \{((pathstart \gamma) \cdot i)..((pathfinish \gamma) \cdot i)\} (\lambda f\text{-}var. (F (f\text{-}var *_R i +$
 $g(f\text{-}var)) \cdot i))$
 $= integral (\{0..1\}) (\lambda x. (F(\gamma(x)) \cdot i) * (vector\text{-}derivative (\lambda x. \gamma(x) \cdot$
 $i) (at x \text{ within } (\{0..1\}))))$

using $gamma\text{-}is\text{-}in\text{-}terms\text{-}of\text{-}i \text{ integral}\text{-}unique[OF \text{ has}\text{-}int] \text{ i}\text{-}norm\text{-}1$

apply ($auto \text{ simp add: pathstart}\text{-}def \text{ pathfinish}\text{-}def$)

apply ($simp \text{ only: real}\text{-}inner\text{-}class.\text{inner}\text{-}add\text{-}left \text{ inner}\text{-}not\text{-}same\text{-}Basis \text{ g}\text{-}orthogonal\text{-}to\text{-}i$
 $\text{ inner}\text{-}scaleR\text{-}left \text{ norm}\text{-}eq\text{-}1$)

by ($auto \text{ simp add: algebra}\text{-}simps$)

have $comp\text{-}in\text{-}eq\text{-}comp\text{-}out: \forall x \in \{0..1\} - s.$
 $(vector\text{-}derivative (\lambda x. \gamma(x) \cdot i) (at x \text{ within } \{0..1\}))$
 $= (vector\text{-}derivative \gamma (at x \text{ within } \{0..1\})) \cdot i$

proof

fix $x:: real$

assume $ass: x \in \{0..1\} - s$

then have $x\text{-}bounds: x \in \{0..1\}$ **by** $auto$

have γ $differentiable \text{ at } x$ **using** $ass \text{ gamma}\text{-}differentiable$ **by** $auto$

then have $dotprod\text{-}in\text{-}is\text{-}out:$
 $vector\text{-}derivative (\lambda x. \gamma(x) \cdot i) (at x)$

= (vector-derivative γ (at x)) \cdot i

using *derivative-component-fun-component*
by force
then have 0 : (vector-derivative γ (at x)) \cdot i = (vector-derivative γ (at x within $\{0..1\}$)) \cdot i
proof –
have (γ has-vector-derivative vector-derivative γ (at x)) (at x)
using $\langle \gamma$ differentiable at $x \rangle$ vector-derivative-works **by blast**
moreover have $0 \leq x \wedge x \leq 1$
using x -bounds **by presburger**
ultimately show ?thesis
by (*simp add: vector-derivative-at-within-ivl*)
qed
have 1 : vector-derivative ($\lambda x. \gamma(x) \cdot i$) (at x) = vector-derivative ($\lambda x. \gamma(x) \cdot i$) (at x within $\{0..1\}$)
using path-vector-deriv-line-integrals **and** vector-derivative-at-within-ivl **and** x -bounds
by (*metis ass atLeastAtMost-iff zero-less-one*)
show vector-derivative ($\lambda x. \gamma x \cdot i$) (at x within $\{0..1\}$) = vector-derivative γ (at x within $\{0..1\}$) \cdot i
using 0 **and** 1 **and** dotprod-in-is-out
by auto
qed
show integral $\{0..1\}$ ($\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i)$) =
integral $\{\text{pathstart } \gamma \cdot i.. \text{pathfinish } \gamma \cdot i\}$ ($\lambda f\text{-var}. F (f\text{-var} *_R i + g f\text{-var}) \cdot i$)
using substitute **and** comp-in-eq-comp-out **and** negligible-finite
Henstock-Kurzweil-Integration.integral-spike
 $[\text{of } s \{0..1\} (\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i))$
 $(\lambda x. F (\gamma x) \cdot i * \text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x \text{ within } \{0..1\}))]$
by (*metis (no-types, lifting) gamma-differentiable(1)*)
have $((\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i))$ has-integral
integral $\{\text{pathstart } \gamma \cdot i.. \text{pathfinish } \gamma \cdot i\}$ ($\lambda f\text{-var}. F (f\text{-var} *_R i + g f\text{-var}) \cdot i$) $\{0..1\}$
using has-int' **and** comp-in-eq-comp-out **and** negligible-finite
Henstock-Kurzweil-Integration.has-integral-spike
 $[\text{of } s \{0..1\} (\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i))$
 $(\lambda x. F (\gamma x) \cdot i * \text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x \text{ within } \{0..1\}))$
integral $\{\text{pathstart } \gamma \cdot i.. \text{pathfinish } \gamma \cdot i\}$ ($\lambda f\text{-var}. F (f\text{-var} *_R i + g f\text{-var})$
 $\cdot i$)]
by (*metis (no-types, lifting) gamma-differentiable(1)*)
then have $(\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i))$
integrable-on $\{0..1\}$
using integrable-on-def **by auto**
then show line-integral-exists $F \{i\} \gamma$
by (*auto simp add: line-integral-exists-def*)
qed

```

lemma line-integral-on-pair-path:
  fixes  $F::('a::\text{euclidean-space}) \Rightarrow ('a)$  and
     $g::\text{real} \Rightarrow 'a$  and
     $\gamma::(\text{real} \Rightarrow 'a)$  and
     $i::'a$ 
  assumes i-norm-1:  $\text{norm } i = 1$  and
    g-orthogonal-to-i:  $\forall x. g(x) \cdot i = 0$  and
    gamma-is-in-terms-of-i:  $\gamma = (\lambda x. f(x) *_{\mathbb{R}} i + g(f(x)))$  and
    gamma-smooth:  $\gamma$  C1-differentiable-on  $\{0 .. 1\}$  and
    g-continuous-on-f: continuous-on  $(f \text{ ' } \{0..1\})$   $g$  and
    path-start-le-path-end:  $(\text{pathstart } \gamma) \cdot i \leq (\text{pathfinish } \gamma) \cdot i$  and
    field-i-comp-cont: continuous-on  $(\text{path-image } \gamma)$   $(\lambda x. F x \cdot i)$ 
  shows  $(\text{line-integral } F \{i\} \gamma)$ 
    =  $\text{integral } (\text{cbox } ((\text{pathstart } \gamma) \cdot i) ((\text{pathfinish } \gamma) \cdot i)) (\lambda f\text{-var}. (F$ 
     $(f\text{-var} *_{\mathbb{R}} i + g(f\text{-var})) \cdot i))$ 
  proof (simp add: line-integral-def)
    have gamma-differentiable:  $\forall x \in \{0 .. 1\}. \gamma$  differentiable at  $x$ 
      using gamma-smooth C1-differentiable-on-eq by blast
    then have gamma-i-component-smooth:
       $\forall x \in \{0 .. 1\}. (\lambda x. \gamma x \cdot i)$  differentiable at  $x$ 
      by auto
    have vec-deriv-i-comp-cont:
      continuous-on  $\{0..1\}$   $(\lambda x. \text{vector-derivative } (\lambda x. \gamma x \cdot i) (\text{at } x \text{ within } \{0..1\}))$ 
    proof –
      have continuous-on  $\{0..1\}$   $(\lambda x. \text{vector-derivative } (\lambda x. \gamma x) (\text{at } x \text{ within } \{0..1\}))$ 
        using gamma-smooth C1-differentiable-on-eq
        by (smt C1-differentiable-on-def atLeastAtMost-iff continuous-on-eq vector-derivative-at-within-ivl)
      then have deriv-comp-cont:
        continuous-on  $\{0..1\}$   $(\lambda x. \text{vector-derivative } (\lambda x. \gamma x) (\text{at } x \text{ within } \{0..1\}) \cdot i)$ 
        by (simp add: continuous-intros)
      show ?thesis
      using derivative-component-fun-component-at-within[OF gamma-differentiable,
of i]
        continuous-on-eq[OF deriv-comp-cont, of  $(\lambda x. \text{vector-derivative } (\lambda x. \gamma x \cdot i)$ 
(at x within  $\{0..1\})$ )]
        by fastforce
    qed
  have field-cont-on-path:
    continuous-on  $((\lambda x. \gamma x \cdot i) \text{ ' } \{0..1\}) (\lambda f\text{-var}. F (f\text{-var} *_{\mathbb{R}} i + g f\text{-var}) \cdot i)$ 
  proof –
    have  $0: (\lambda x. \gamma x \cdot i) = f$ 
    proof
      fix  $x$ 
      show  $\gamma x \cdot i = f x$ 
      using g-orthogonal-to-i i-norm-1
      by (simp only: gamma-is-in-terms-of-i real-inner-class.inner-add-left g-orthogonal-to-i inner-scaleR-left inner-same-Basis norm-eq-1)
    qed
  qed

```

```

show ?thesis
unfolding 0
apply (rule continuous-on-compose2 [of - ( $\lambda x. F(x) \cdot i$ ) f ' { 0..1} ( $\lambda x. x$ 
*_R i + g x)]
      field-i-comp-cont g-continuous-on-f field-i-comp-cont continuous-intros)+
by (auto simp add: gamma-is-in-terms-of-i path-image-def)
qed
have path-vector-deriv-line-integrals:
   $\forall x \in \{0..1\}. ((\lambda x. \gamma x \cdot i)$  has-vector-derivative
      vector-derivative ( $\lambda x. \gamma x \cdot i$ ) (at x))
      (at x)
using gamma-i-component-smooth and derivative-component-fun-component
and
  vector-derivative-works
by blast
then have  $\forall x \in \{0..1\}. ((\lambda x. \gamma x \cdot i)$  has-vector-derivative
      vector-derivative ( $\lambda x. \gamma x \cdot i$ ) (at x within {0..1}))
      (at x within {0..1})
using has-vector-derivative-at-within vector-derivative-at-within-ivl
by fastforce
then have substitute:
  integral (cbox ((pathstart  $\gamma$ )  $\cdot i$ ) ((pathfinish  $\gamma$ )  $\cdot i$ )) ( $\lambda f$ -var. (F (f-var *_R i +
g(f-var))  $\cdot i$ ))
    = integral (cbox 0 1) ( $\lambda x. (F(\gamma(x)) \cdot i) * (\text{vector-derivative } (\lambda x. \gamma(x))$ 
 $\cdot i)$  (at x within {0..1})))
using gauge-integral-by-substitution
  [of 0 1 ( $\lambda x. (\gamma x) \cdot i$ )
    ( $\lambda x. \text{vector-derivative } (\lambda x. \gamma(x)) \cdot i$ ) (at x within {0..1})
    ( $\lambda f$ -var. (F (f-var *_R i + g(f-var))  $\cdot i$ ))] and
  path-start-le-path-end and vec-deriv-i-comp-cont and field-cont-on-path and
  gamma-is-in-terms-of-i i-norm-1
apply (auto simp add: pathstart-def pathfinish-def)
apply (simp only: real-inner-class.inner-add-left inner-not-same-Basis g-orthogonal-to-i
inner-scaleR-left norm-eq-1)
by (auto)

have comp-in-eq-comp-out:  $\forall x \in \{0..1\}.
  (\text{vector-derivative } (\lambda x. \gamma(x)) \cdot i)$  (at x within {0..1})
    = ( $\text{vector-derivative } \gamma$  (at x within {0..1}))  $\cdot i$ 
proof
fix x:: real
assume x-bounds:  $x \in \{0..1\}$ 
then have  $\gamma$  differentiable at x using gamma-differentiable by auto
then have dotprod-in-is-out:
  vector-derivative ( $\lambda x. \gamma(x) \cdot i$ ) (at x)
    = ( $\text{vector-derivative } \gamma$  (at x))  $\cdot i$ 
using derivative-component-fun-component
by force
then have 0: ( $\text{vector-derivative } \gamma$  (at x))  $\cdot i$ 

```

= (vector-derivative γ (at x within $\{0..1\}$)) $\cdot i$

using *has-vector-derivative-at-within* **and** *x-bounds* **and** *vector-derivative-at-within-ivl*
by (*smt atLeastAtMost-iff gamma-differentiable inner-commute vector-derivative-works*)
have 1: vector-derivative ($\lambda x. \gamma(x) \cdot i$) (at x) = vector-derivative ($\lambda x. \gamma(x) \cdot$
i) (at x within $\{0..1\}$)
using *path-vector-deriv-line-integrals* **and** *vector-derivative-at-within-ivl* **and**
x-bounds
by *fastforce*
show vector-derivative ($\lambda x. \gamma x \cdot i$) (at x within $\{0..1\}$) = vector-derivative γ
(at x within $\{0..1\}$) $\cdot i$
using 0 **and** 1 **and** *dotprod-in-is-out*
by *auto*
qed
show integral $\{0..1\}$ ($\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\})$
 $\cdot i)$) =
integral $\{\text{pathstart } \gamma \cdot i.. \text{pathfinish } \gamma \cdot i\}$ ($\lambda f\text{-var. } F (f\text{-var} *_{\mathbb{R}} i +$
 $g f\text{-var}) \cdot i$)
using *substitute* **and** *comp-in-eq-comp-out* **and**
Henstock-Kurzweil-Integration.integral-cong
[*of* $\{0..1\}$ ($\lambda x. F (\gamma x) \cdot i * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot i$)
($\lambda x. F (\gamma x) \cdot i * \text{vector-derivative } (\lambda x. \gamma x \cdot i) \text{ (at } x \text{ within } \{0..1\})$)]
by *auto*
qed

lemma *content-box-cases*:

content (box a b) = (if $\forall i \in \text{Basis. } a \cdot i \leq b \cdot i$ then prod ($\lambda i. b \cdot i - a \cdot i$) Basis else
0)
by (*simp add: measure-lborel-box-eq inner-diff*)

lemma *content-box-cbox*:

shows *content* (box a b) = *content* (cbox a b)
by (*simp add: content-box-cases content-cbox-cases*)

lemma *content-eq-0*: *content* (box a b) = 0 \iff ($\exists i \in \text{Basis. } b \cdot i \leq a \cdot i$)

by (*auto simp: content-box-cases not-le intro: less-imp-le antisym eq-refl*)

lemma *content-pos-lt-eq*: $0 < \text{content} (\text{cbox } a (b::'a::\text{euclidean-space})) \iff (\forall i \in \text{Basis. } a \cdot i < b \cdot i)$

by (*auto simp add: content-cbox-cases less-le prod-nonneg*)

lemma *content-lt-nz*: $0 < \text{content} (\text{box } a \ b) \iff \text{content} (\text{box } a \ b) \neq 0$

using *Paths.content-eq-0 zero-less-measure-iff* **by** *blast*

lemma *content-subset*: $\text{cbox } a \ b \subseteq \text{box } c \ d \implies \text{content} (\text{cbox } a \ b) \leq \text{content} (\text{box } c \ d)$

unfolding *measure-def*

by (*intro enn2real-mono emeasure-mono*) (*auto simp: emeasure-lborel-cbox-eq emeasure-lborel-box-eq*)

lemma *sum-content-null*:
assumes *content (box a b) = 0*
and *p tagged-division-of (box a b)*
shows *sum (λ(x,k). content k *_R f x) p = (0::'a::real-normed-vector)*
proof (*rule sum.neutral, rule*)
fix *y*
assume *y: y ∈ p*
obtain *x k where xk: y = (x, k)*
using *surj-pair[of y] by blast*
note *assm = tagged-division-ofD(3-4)[OF assms(2) y[unfolded xk]]*
from *this(2) obtain c d where k: k = cbox c d by blast*
have *(λ(x, k). content k *_R f x) y = content k *_R f x*
unfolding *xk by auto*
also have *... = 0*
using *content-subset[OF assm(1)[unfolded k]] content-pos-le[of cbox c d]*
unfolding *assms(1) k*
by *auto*
finally show *(λ(x, k). content k *_R f x) y = 0 .*
qed

lemma *has-integral-null [intro]: content(box a b) = 0 ⇒ (f has-integral 0) (box a b)*
by (*simp add: content-box-cbox content-eq-0-interior*)

lemma *line-integral-distrib*:
assumes *line-integral-exists f basis g1*
line-integral-exists f basis g2
valid-path g1 valid-path g2
shows *line-integral f basis (g1 +++ g2) = line-integral f basis g1 + line-integral f basis g2*
line-integral-exists f basis (g1 +++ g2)
proof –
obtain *s1 s2*
where *s1: finite s1 ∀ x∈{0..1} – s1. g1 differentiable at x*
and *s2: finite s2 ∀ x∈{0..1} – s2. g2 differentiable at x*
using *assms*
by (*auto simp: valid-path-def piecewise-C1-differentiable-on-def C1-differentiable-on-eq*)
obtain *i1 i2*
where *1: ((λx. ∑ b∈basis. f (g1 x) · b * (vector-derivative g1 (at x within {0..1}) · b)) has-integral i1) {0..1}*
and *2: ((λx. ∑ b∈basis. f (g2 x) · b * (vector-derivative g2 (at x within {0..1}) · b)) has-integral i2) {0..1}*
using *assms*
by (*auto simp: line-integral-exists-def*)
have *i1: ((λx. 2 * (∑ b∈basis. f (g1 (2 * x)) · b * (vector-derivative g1 (at (2 * x) within {0..1}) · b))) has-integral i1) {0..1/2}*
and *i2: ((λx. 2 * (∑ b∈basis. f (g2 (2 * x - 1)) · b * (vector-derivative g2 (at ((2 * x) - 1) within {0..1}) · b))) has-integral i2) {1/2..1}*

```

using has-integral-affinity01 [OF 1, where  $m=2$  and  $c=0$ , THEN has-integral-cmul
[where  $c=2$ ]]
  has-integral-affinity01 [OF 2, where  $m=2$  and  $c=-1$ , THEN has-integral-cmul
[where  $c=2$ ]]
  by (simp-all only: image-affinity-atLeastAtMost-div-diff, simp-all add: scaleR-conv-of-real
mult-ac)
  have g1:  $\llbracket 0 \leq z; z*2 < 1; z*2 \notin s1 \rrbracket \implies$ 
    vector-derivative  $(\lambda x. \text{if } x*2 \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1))$  (at z
within  $\{0..1\}$ ) =
    2 *R vector-derivative g1 (at (z*2) within  $\{0..1\}$ ) for z
  proof -
    have i: $\llbracket 0 \leq z; z*2 < 1; z*2 \notin s1 \rrbracket \implies$ 
      vector-derivative  $(\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1))$ 
(at z within  $\{0..1\}$ ) = 2 *R vector-derivative g1 (at (z * 2)) for z
    proof -
      have g1-at-z: $\llbracket 0 \leq z; z*2 < 1; z*2 \notin s1 \rrbracket \implies$ 
         $((\lambda x. \text{if } x*2 \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1))$ 
has-vector-derivative
          2 *R vector-derivative g1 (at (z*2))) (at z) for z
      apply (rule has-vector-derivative-transform-at [of  $|z - 1/2| - (\lambda x. g1(2*x))$ ])
      apply (simp-all add: dist-real-def abs-if split: if-split-asm)
      apply (rule vector-diff-chain-at [of  $\lambda x. 2*x \ 2 - g1$ , simplified o-def])
      apply (simp add: has-vector-derivative-def has-derivative-def bounded-linear-mult-left)
      using s1
      apply (auto simp: algebra-simps vector-derivative-works)
      done
      assume ass:  $0 \leq z \ z*2 < 1 \ z*2 \notin s1$ 
      then have z-ge:  $z \leq 1$  by auto
      show vector-derivative  $(\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x -$ 
1)) (at z within  $\{0..1\}$ ) = 2 *R vector-derivative g1 (at (z * 2))
      using Derivative.vector-derivative-at-within-ivl[OF g1-at-z[OF ass] ass(1)
z-ge]
      by auto
    qed
    assume ass:  $0 \leq z \ z*2 < 1 \ z*2 \notin s1$ 
    then have (g1 has-vector-derivative ((vector-derivative g1 (at (z*2)))) (at
(z*2)))
      using s1 by (auto simp: algebra-simps vector-derivative-works)
    then have ii: (vector-derivative g1 (at (z*2) within  $\{0..1\}$ )) = (vector-derivative
g1 (at (z*2)))
      using Derivative.vector-derivative-at-within-ivl ass
      by force
    show vector-derivative  $(\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1))$ 
(at z within  $\{0..1\}$ ) = 2 *R vector-derivative g1 (at (z * 2) within  $\{0..1\}$ )
      using i[OF ass] ii
      by auto
    qed
  have g2:  $\llbracket 1 < z*2; z \leq 1; z*2 - 1 \notin s2 \rrbracket \implies$ 
    vector-derivative  $(\lambda x. \text{if } x*2 \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1))$  (at z

```

$within \{0..1\} =$
 $2 *_{\mathbb{R}} \text{vector-derivative } g2 \text{ (at } (z*2 - 1) \text{ within } \{0..1\}) \text{ for } z \quad \text{proof}$
 $-$
have $i: [1 < z*2; z \leq 1; z*2 - 1 \notin s2] \implies$
 $\text{vector-derivative } (\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1))$
 $(\text{at } z \text{ within } \{0..1\})$
 $= 2 *_{\mathbb{R}} \text{vector-derivative } g2 \text{ (at } (z * 2 - 1)) \text{ for } z$
proof-
have $g2\text{-at-}z: [1 < z*2; z \leq 1; z*2 - 1 \notin s2] \implies$
 $((\lambda x. \text{if } x*2 \leq 1 \text{ then } g1 (2*x) \text{ else } g2 (2*x - 1))$
 $\text{has-vector-derivative } 2 *_{\mathbb{R}} \text{vector-derivative } g2 \text{ (at } (z*2 - 1))) \text{ (at } z) \text{ for } z$
apply $(\text{rule has-vector-derivative-transform-at [of } |z - 1/2| - (\lambda x. g2 (2*x$
 $- 1))])$
apply $(\text{simp-all add: dist-real-def abs-if split: if-split-asm})$
apply $(\text{rule vector-diff-chain-at [of } \lambda x. 2*x - 1 \text{ - } g2, \text{ simplified o-def]})$
apply $(\text{simp add: has-vector-derivative-def has-derivative-def bounded-linear-mult-left})$
using $s2$
apply $(\text{auto simp: algebra-simps vector-derivative-works})$
done
assume $\text{ass: } 1 < z*2 \text{ } z \leq 1 \text{ } z*2 - 1 \notin s2$
then have $z\text{-le: } z \leq 1 \text{ by auto}$
have $z\text{-ge: } 0 \leq z \text{ using ass by auto}$
show $\text{vector-derivative } (\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x -$
 $1)) \text{ (at } z \text{ within } \{0..1\})$
 $= 2 *_{\mathbb{R}} \text{vector-derivative } g2 \text{ (at } (z * 2 - 1))$
using $\text{Derivative.vector-derivative-at-within-ivl [OF } g2\text{-at-}z[\text{OF ass}] \text{ } z\text{-ge } z\text{-le}]$
by auto
qed
assume $\text{ass: } 1 < z*2 \text{ } z \leq 1 \text{ } z*2 - 1 \notin s2$
then have $(g2 \text{ has-vector-derivative } ((\text{vector-derivative } g2 \text{ (at } (z*2 - 1))))))$
 $(\text{at } (z*2 - 1))$
using $s2 \text{ by (auto simp: algebra-simps vector-derivative-works)}$
then have $ii: (\text{vector-derivative } g2 \text{ (at } (z*2 - 1) \text{ within } \{0..1\})) = (\text{vector-derivative}$
 $g2 \text{ (at } (z*2 - 1)))$
using $\text{Derivative.vector-derivative-at-within-ivl ass}$
by force
show $\text{vector-derivative } (\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1))$
 $(\text{at } z \text{ within } \{0..1\}) = 2 *_{\mathbb{R}} \text{vector-derivative } g2 \text{ (at } (z * 2 - 1) \text{ within } \{0..1\})$
using $i[\text{OF ass}] \text{ } ii$
by auto
qed
have $\text{lem1: } ((\lambda x. \sum_{b \in \text{basis}} f ((g1+++g2) x) \cdot b * (\text{vector-derivative } (g1+++g2)$
 $(\text{at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } i1 \{0..1/2\}$
apply $(\text{rule has-integral-spike-finite [OF - - } i1, \text{ of insert } (1/2) ((*)2 - 's1)])$
using $s1$
apply $(\text{force intro: finite-vimageI [where } h = (*)2 \text{] inj-onI})$
apply $(\text{clarsimp simp add: joinpaths-def scaleR-conv-of-real mult-ac } g1)$
by $(\text{simp add: sum-distrib-left})$
moreover have $\text{lem2: } ((\lambda x. \sum_{b \in \text{basis}} f ((g1+++g2) x) \cdot b * (\text{vector-derivative}$

```

(g1+++g2) (at x within {0..1}) · b)) has-integral i2) {1/2..1}
  apply (rule has-integral-spike-finite [OF - - i2, of insert (1/2) ((λx. 2*x-1)
- ' s2)])
  using s2
  apply (force intro: finite-vimageI [where h = λx. 2*x-1] inj-onI)
  apply (clarsimp simp add: joinpaths-def scaleR-conv-of-real mult-ac g2)
  by (simp add: sum-distrib-left)
ultimately
show line-integral f basis (g1 +++ g2) = line-integral f basis g1 + line-integral
f basis g2
  apply (simp add: line-integral-def)
  apply (rule integral-unique [OF has-integral-combine [where c = 1/2]])
  using 1 2 integral-unique apply auto
done
show line-integral-exists f basis (g1 +++ g2)
proof (simp add: line-integral-exists-def integrable-on-def)
  have (1::real) ≤ 1 * 2 ∧ (0::real) ≤ 1 / 2
  by simp
  then show ∃ r. ((λr. ∑ a∈basis. f ((g1 +++ g2) r) · a * (vector-derivative
(g1 +++ g2) (at r within {0..1}) · a)) has-integral r) {0..1}
  using has-integral-combine [where c = 1/2] 1 2 divide-le-eq-numeral1(1) lem1
lem2 by blast
qed
qed

```

lemma *line-integral-exists-joinD1*:

```

assumes line-integral-exists f basis (g1 +++ g2) valid-path g1
shows line-integral-exists f basis g1
proof -
  obtain s1
  where s1: finite s1 ∀ x∈{0..1} - s1. g1 differentiable at x
  using assms by (auto simp: valid-path-def piecewise-C1-differentiable-on-def
C1-differentiable-on-eq)
  have (λx. ∑ b∈basis. f ((g1 +++ g2) (x/2)) · b * (vector-derivative (g1 +++
g2) (at (x/2) within {0..1}) · b)) integrable-on {0..1}
  using assms
  apply (auto simp: line-integral-exists-def)
  apply (drule integrable-on-subbox [where a=0 and b=1/2])
  apply (auto intro: integrable-affinity [of - 0 1/2::real 1/2 0, simplified])
done
then have *: (λx. ∑ b∈basis. ((f ((g1 +++ g2) (x/2)) · b) / 2) * (vector-derivative
(g1 +++ g2) (at (x/2) within {0..1}) · b)) integrable-on {0..1}
  by (auto simp: Groups-Big.sum-distrib-left dest: integrable-cmul [where c=1/2]
simp: scaleR-conv-of-real)
  have g1: [0 ≤ z; z*2 < 1; z*2 ∉ s1] ⇒
    vector-derivative (λx. if x*2 ≤ 1 then g1 (2*x) else g2 (2*x - 1)) (at z
within {0..1}) =
    2 *_R vector-derivative g1 (at (z*2) within {0..1}) for z
proof -

```



```

have i:[ $0 \leq z; z * 2 < 1; z * 2 \notin s1$ ]  $\implies$ 
  vector-derivative ( $\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
(at z within {0..1}) = 2 *R vector-derivative g1 (at (z * 2)) for z
proof –
  have g1-at-z:[ $0 \leq z; z * 2 < 1; z * 2 \notin s1$ ]  $\implies$ 
    (( $\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
has-vector-derivative
  2 *R vector-derivative g1 (at (z * 2))) (at z) for z
apply (rule has-vector-derivative-transform-at [of |z - 1/2| - ( $\lambda x. g1(2 * x)$ )])
  apply (simp-all add: dist-real-def abs-if split: if-split-asm)
  apply (rule vector-diff-chain-at [of  $\lambda x. 2 * x * 2 - g1$ , simplified o-def])
apply (simp add: has-vector-derivative-def has-derivative-def bounded-linear-mult-left)
  using s1
  apply (auto simp: algebra-simps vector-derivative-works)
  done
assume ass:  $0 \leq z \wedge z * 2 < 1 \wedge z * 2 \notin s1$ 
then have z-ge:  $z \leq 1$  by auto
  show vector-derivative ( $\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
(at z within {0..1}) = 2 *R vector-derivative g1 (at (z * 2))
  using Derivative.vector-derivative-at-within-ivl[OF g1-at-z[OF ass] ass(1)
z-ge]
  by auto
qed
assume ass:  $0 \leq z \wedge z * 2 < 1 \wedge z * 2 \notin s1$ 
then have (g1 has-vector-derivative ((vector-derivative g1 (at (z * 2)))) (at
(z * 2)))
  using s1 by (auto simp: algebra-simps vector-derivative-works)
then have ii: (vector-derivative g1 (at (z * 2) within {0..1})) = (vector-derivative
g1 (at (z * 2)))
  using Derivative.vector-derivative-at-within-ivl ass by force
  show vector-derivative ( $\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)$ )
(at z within {0..1}) = 2 *R vector-derivative g1 (at (z * 2) within {0..1})
  using i[OF ass] ii by auto
qed
show ?thesis
  using s1
  apply (auto simp: line-integral-exists-def)
  apply (rule integrable-spike-finite [of {0,1}  $\cup$  s1, OF - - *])
  apply (auto simp: joinpaths-def scaleR-conv-of-real g1)
  done
qed

```

lemma line-integral-exists-joinD2:

assumes line-integral-exists f basis (g1 +++ g2) valid-path g2
shows line-integral-exists f basis g2

proof –

obtain s2

where s2: finite s2 $\forall x \in \{0..1\} - s2. g2$ differentiable at x

using assms **by** (auto simp: valid-path-def piecewise-C1-differentiable-on-def

C1-differentiable-on-eq

```

have ( $\lambda x. \sum b \in \text{basis}. f ((g1 \text{ +++ } g2) (x/2 + 1/2)) \cdot b * (\text{vector-derivative } (g1 \text{ +++ } g2) (\text{at } (x/2 + 1/2) \text{ within } \{0..1\}) \cdot b)) \text{ integrable-on } \{0..1\}$ 
using assms
apply (auto simp: line-integral-exists-def)
apply (drule integrable-on-subcbox [where a=1/2 and b=1], auto)
apply (drule integrable-affinity [of - 1/2::real 1 1/2 1/2, simplified])
apply (simp add: image-affinity-atLeastAtMost-diff)
done
then have  $*$ : ( $\lambda x. \sum b \in \text{basis}. ((f ((g1 \text{ +++ } g2) (x/2 + 1/2)) \cdot b) / 2) * (\text{vector-derivative } (g1 \text{ +++ } g2) (\text{at } (x/2 + 1/2) \text{ within } \{0..1\}) \cdot b)) \text{ integrable-on } \{0..1\}$ 
by (auto simp: Groups-Big.sum-distrib-left dest: integrable-cmul [where c=1/2] simp: scaleR-conv-of-real)
have  $g2$ : [ $1 < z * 2; z \leq 1; z * 2 - 1 \notin s2$ ]  $\implies$ 
 $\text{vector-derivative } (\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)) (\text{at } z \text{ within } \{0..1\}) =$ 
 $2 *_{\mathbb{R}} \text{vector-derivative } g2 (\text{at } (z * 2 - 1) \text{ within } \{0..1\})$  for  $z$  proof
-
have  $i$ : [ $1 < z * 2; z \leq 1; z * 2 - 1 \notin s2$ ]  $\implies$ 
 $\text{vector-derivative } (\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)) (\text{at } z \text{ within } \{0..1\}) =$ 
 $2 *_{\mathbb{R}} \text{vector-derivative } g2 (\text{at } (z * 2 - 1))$  for  $z$ 
proof -
have  $g2\text{-at-}z$ : [ $1 < z * 2; z \leq 1; z * 2 - 1 \notin s2$ ]  $\implies$ 
 $((\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1))) (\text{at } z) \text{ for } z$ 
 $\text{has-vector-derivative } 2 *_{\mathbb{R}} \text{vector-derivative } g2 (\text{at } (z * 2 - 1))$ 
apply (rule has-vector-derivative-transform-at [of |z - 1/2| - ( $\lambda x. g2 (2 * x - 1)$ )])
apply (simp-all add: dist-real-def abs-if split: if-split-asm)
apply (rule vector-diff-chain-at [of  $\lambda x. 2 * x - 1$  2 -  $g2$ , simplified o-def])
apply (simp add: has-vector-derivative-def has-derivative-def bounded-linear-mult-left)
using  $s2$ 
apply (auto simp: algebra-simps vector-derivative-works)
done
assume  $ass$ :  $1 < z * 2$   $z \leq 1$   $z * 2 - 1 \notin s2$ 
then have  $z\text{-le}$ :  $z \leq 1$  by auto
have  $z\text{-ge}$ :  $0 \leq z$  using  $ass$  by auto
show  $\text{vector-derivative } (\lambda x. \text{if } x * 2 \leq 1 \text{ then } g1 (2 * x) \text{ else } g2 (2 * x - 1)) (\text{at } z \text{ within } \{0..1\}) =$ 
 $2 *_{\mathbb{R}} \text{vector-derivative } g2 (\text{at } (z * 2 - 1))$ 
using Derivative.vector-derivative-at-within-ivl[OF  $g2\text{-at-}z$ [OF  $ass$ ]  $z\text{-ge}$   $z\text{-le}$ ]
by auto
qed
assume  $ass$ :  $1 < z * 2$   $z \leq 1$   $z * 2 - 1 \notin s2$ 
then have ( $g2 \text{ has-vector-derivative } ((\text{vector-derivative } g2 (\text{at } (z * 2 - 1))))$ )
( $\text{at } (z * 2 - 1)$ )
using  $s2$  by (auto simp: algebra-simps vector-derivative-works)
then have  $ii$ : ( $\text{vector-derivative } g2 (\text{at } (z * 2 - 1) \text{ within } \{0..1\}) = (\text{vector-derivative$ 

```

```

g2 (at (z*2 - 1)))
  using Derivative.vector-derivative-at-within-ivl ass
  by force
  show vector-derivative (λx. if x * 2 ≤ 1 then g1 (2 * x) else g2 (2 * x - 1))
(at z within {0..1}) = 2 *R vector-derivative g2 (at (z * 2 - 1) within {0..1})
  using i[OF ass] ii
  by auto
qed
show ?thesis
  using s2
  apply (auto simp: line-integral-exists-def)
  apply (rule integrable-spike-finite [of {0,1} ∪ s2, OF - - *])
  apply (auto simp: joinpaths-def scaleR-conv-of-real g2)
done
qed

```

lemma *has-line-integral-on-reverse-path:*

assumes *g: valid-path g and int:*

$((\lambda x. \sum_{b \in \text{basis}. F} (g \ x) \cdot b * (\text{vector-derivative } g \ (\text{at } x \ \text{within } \{0..1\}) \cdot b))$
has-integral c){0..1}

shows $((\lambda x. \sum_{b \in \text{basis}. F} ((\text{reversepath } g) \ x) \cdot b * (\text{vector-derivative } (\text{reversepath } g) \ (\text{at } x \ \text{within } \{0..1\}) \cdot b))$
has-integral -c){0..1}

proof –

{ **fix** *s x*

assume *xs: g C1-differentiable-on ({0..1} - s) x ∉ (-) 1 ‘ s 0 ≤ x x ≤ 1*

have $\text{vector-derivative } (\lambda x. g \ (1 - x)) \ (\text{at } x \ \text{within } \{0..1\}) =$
 $-\ \text{vector-derivative } g \ (\text{at } (1 - x) \ \text{within } \{0..1\})$

proof –

obtain *f' where f': (g has-vector-derivative f') (at (1 - x))*

using *xs*

by (*force simp: has-vector-derivative-def C1-differentiable-on-def*)

have $(g \ o \ (\lambda x. 1 - x) \ \text{has-vector-derivative } -1 *_R f') \ (\text{at } x)$

apply (*rule vector-diff-chain-within*)

apply (*intro vector-diff-chain-within derivative-eq-intros | simp*)+

apply (*rule has-vector-derivative-at-within [OF f']*)

done

then have $m f': ((\lambda x. g \ (1 - x)) \ \text{has-vector-derivative } -f') \ (\text{at } x)$

by (*simp add: o-def*)

show *?thesis*

using *xs*

by (*auto simp: vector-derivative-at-within-ivl [OF mf'] vector-derivative-at-within-ivl*
 $[OF f']$)

qed

} **note** $* = \text{this}$

obtain *S where continuous-on {0..1} g finite S g C1-differentiable-on {0..1}*
 $- S$

using *g*

by (*auto simp: valid-path-def piecewise-C1-differentiable-on-def*)

then show *?thesis*

using *has-integral-affinity01* [*OF int*, **where** $m = -1$ **and** $c = 1$]
unfolding *reversepath-def*
by (*rule-tac* $S = (\lambda x. 1 - x)$ ‘*S in has-integral-spike-finite*) (*auto simp*: *
has-integral-neg Groups-Big.sum-negf)
qed

lemma *line-integral-on-reverse-path*:

assumes *valid-path* γ *line-integral-exists* F *basis* γ
shows *line-integral* F *basis* $\gamma = -$ (*line-integral* F *basis* (*reversepath* γ))
line-integral-exists F *basis* (*reversepath* γ)
proof –
obtain i **where**
 $0: ((\lambda x. \sum_{b \in \text{basis}} F (\gamma x) \cdot b * (\text{vector-derivative } \gamma \text{ (at } x \text{ within } \{0..1\}) \cdot b))$
has-integral i) $\{0..1\}$
using *assms unfolding integrable-on-def line-integral-exists-def* **by** *auto*
then have $1: ((\lambda x. \sum_{b \in \text{basis}} F ((\text{reversepath } \gamma) x) \cdot b * (\text{vector-derivative}$
(*reversepath* γ) (at x within $\{0..1\}$) $\cdot b))$ *has-integral* $-i$) $\{0..1\}$
using *has-line-integral-on-reverse-path assms*
by *auto*
then have *rev-line-integral:line-integral* F *basis* (*reversepath* γ) = $-i$
using *line-integral-def Henstock-Kurzweil-Integration.integral-unique*
by (*metis (no-types)*)
have *line-integral: line-integral* F *basis* $\gamma = i$
using *line-integral-def 0 Henstock-Kurzweil-Integration.integral-unique*
by *blast*
show *line-integral* F *basis* $\gamma = -$ (*line-integral* F *basis* (*reversepath* γ))
using *line-integral rev-line-integral*
by *auto*
show *line-integral-exists* F *basis* (*reversepath* γ)
using 1 *line-integral-exists-def*
by *auto*
qed

lemma *line-integral-exists-on-degenerate-path*:

assumes *finite basis*
shows *line-integral-exists* F *basis* $(\lambda x. c)$
proof –
have *every-component-integrable*:
 $\forall b \in \text{basis}. (\lambda x. F ((\lambda x. c) x) \cdot b * (\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\})$
 $\cdot b))$ *integrable-on* $\{0..1\}$
proof
fix b
assume *b-in-basis*: $b \in \text{basis}$
have *cont-field-zero-one: continuous-on* $\{0..1\}$ $(\lambda x. F ((\lambda x. c) x) \cdot b)$
using *continuous-on-const by fastforce*
have *cont-path-zero-one*:
continuous-on $\{0..1\}$ $(\lambda x. (\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\})) \cdot b)$
proof –
have $((\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\})) \cdot b) = 0$ **if** $x \in \{0..1\}$

for x
proof –
have $\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) = 0$
using $\text{that gamma-deriv-at-within[of } 0 \ 1] \text{ differentiable-const vector-derivative-const-at}$
by fastforce
then show $\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) \cdot b = 0$
by auto
qed
then show $\text{continuous-on } \{0..1\} (\lambda x. (\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\})) \cdot b)$
using $\text{continuous-on-const[of } \{0..1\} \ 0] \text{ continuous-on-eq[of } \{0..1\} \ \lambda x. \ 0$
 $(\lambda x. (\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\})) \cdot b)$
by auto
qed
show $(\lambda x. F \ (c) \cdot b * (\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) \cdot b))$
 $\text{integrable-on } \{0..1\}$
using $\text{cont-field-zero-one cont-path-zero-one continuous-on-mult integrable-continuous-real}$
by blast
qed
have $\text{integrable-sum': } \bigwedge t \ f \ s. \ \text{finite } t \implies$
 $\forall a \in t. \ f \ a \ \text{integrable-on } s \implies (\lambda x. \sum a \in t. \ f \ a \ x) \ \text{integrable-on } s$
using $\text{integrable-sum by metis}$
have $\text{field-integrable-on-basis:}$
 $(\lambda x. \sum b \in \text{basis}. \ F \ (c) \cdot b * (\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) \cdot b))$
 $\text{integrable-on } \{0..1\}$
using $\text{integrable-sum'[OF assms(1) every-component-integrable]}$
by auto
then show ?thesis
using $\text{line-integral-exists-def by auto}$
qed

lemma $\text{degenerate-path-is-valid-path: valid-path } (\lambda x. c)$
by $(\text{auto simp add: valid-path-def piecewise-C1-differentiable-on-def continuous-on-const})$

lemma $\text{line-integral-degenerate-path:}$
assumes finite basis
shows $\text{line-integral } F \ \text{basis } (\lambda x. c) = 0$
proof $(\text{simp add: line-integral-def})$
have $((\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\})) \cdot b) = 0$ **if** $x \in \{0..1\}$ **for**
 $x \ b$
proof –
have $\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) = 0$
using $\text{that gamma-deriv-at-within[of } 0 \ 1] \text{ differentiable-const vector-derivative-const-at}$
by fastforce
then show $\text{vector-derivative } (\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) \cdot b = 0$
by auto
qed
then have $0: \bigwedge x. \ x \in \{0..1\} \implies (\lambda x. \sum b \in \text{basis}. \ F \ c \cdot b * (\text{vector-derivative}$
 $(\lambda x. c) \text{ (at } x \text{ within } \{0..1\}) \cdot b)) \ x = (\lambda x. \ 0) \ x$

by auto
 then show $\text{integral } \{0..1\} (\lambda x. \sum_{b \in \text{basis}. F c \cdot b * (\text{vector-derivative } (\lambda x. c) (\text{at } x \text{ within } \{0..1\}) \cdot b)) = 0$
 using $\text{integral-cong}[\text{of } \{0..1\}, OF 0]$ integral-0 by auto
 qed

definition *point-path* where
 $\text{point-path } \gamma \equiv \exists c. \gamma = (\lambda x. c)$

lemma *line-integral-point-path*:
 assumes *point-path* γ
 assumes *finite basis*
 shows $\text{line-integral } F \text{ basis } \gamma = 0$
 using $\text{assms}(1)$ *point-path-def* $\text{line-integral-degenerate-path}[OF \text{assms}(2)]$
 by force

lemma *line-integral-exists-point-path*:
 assumes *finite basis* *point-path* γ
 shows $\text{line-integral-exists } F \text{ basis } \gamma$
 using assms
 apply(*simp add: point-path-def*)
 using $\text{line-integral-exists-on-degenerate-path}$ by auto

lemma *line-integral-exists-subpath*:
 assumes f : $\text{line-integral-exists } f \text{ basis } g$ and g : *valid-path* g
 and uv : $u \in \{0..1\} v \in \{0..1\} u \leq v$
 shows $(\text{line-integral-exists } f \text{ basis } (\text{subpath } u v g))$
proof (*cases* $v=u$)
 case *tr*: *True*
 have zero : $(\sum_{b \in \text{basis}. f (g u) \cdot b * (\text{vector-derivative } (\lambda x. g u) (\text{at } x \text{ within } \{0..1\}) \cdot b)) = 0$ if $x \in \{0..1\}$ for $x::\text{real}$
proof –
 have $(\text{vector-derivative } (\lambda x. g u) (\text{at } x \text{ within } \{0..1\})) = 0$
 using $\text{Deriv.has-vector-derivative-const}$ that $\text{Derivative.vector-derivative-at-within-ivl}$
 by *fastforce*
 then show $(\sum_{b \in \text{basis}. f (g u) \cdot b * (\text{vector-derivative } (\lambda x. g u) (\text{at } x \text{ within } \{0..1\}) \cdot b)) = 0$
 by auto
 qed
 then have $(\lambda x. \sum_{b \in \text{basis}. f (g u) \cdot b * (\text{vector-derivative } (\lambda x. g u) (\text{at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } 0$ $\{0..1\}$
 by (*meson has-integral-is-0*)
 then show *?thesis*
 using f *tr* by (*auto simp add: line-integral-def line-integral-exists-def sub-path-def*)
next
 case *False*
 obtain s where s : $\bigwedge x. x \in \{0..1\} - s \implies g$ *differentiable at* x and fs : *finite* s
 using g *unfolding* $\text{piecewise-C1-differentiable-on-def}$ $\text{C1-differentiable-on-eq}$

valid-path-def **by** *blast*
have *: $((\lambda x. \sum_{b \in \text{basis}} f (g ((v - u) * x + u)) \cdot b * (\text{vector-derivative } g \text{ (at } ((v - u) * x + u) \text{ within } \{0..1\}) \cdot b))$
 $\text{has-integral } (1 / (v - u)) * \text{integral } \{u..v\} (\lambda x. \sum_{b \in \text{basis}} f (g (x)) \cdot b * (\text{vector-derivative } g \text{ (at } x \text{ within } \{0..1\}) \cdot b)))$
 $\{0..1\}$
using *f uv*
apply (*simp add: line-integral-exists-def subpath-def*)
apply (*drule integrable-on-subcbox [where a=u and b=v, simplified]*)
apply (*simp-all add: has-integral-integral*)
apply (*drule has-integral-affinity [where m=v-u and c=u, simplified]*)
apply (*simp-all add: False image-affinity-atLeastAtMost-div-diff scaleR-conv-of-real*)
apply (*simp add: divide-simps False*)
done
have $vd: \bigwedge x. x \in \{0..1\} \implies$
 $x \notin (\lambda t. (v - u) *_R t + u) - ' s \implies$
 $\text{vector-derivative } (\lambda x. g ((v - u) * x + u)) \text{ (at } x \text{ within } \{0..1\}) = (v - u)$
 $*_R \text{vector-derivative } g \text{ (at } ((v - u) * x + u) \text{ within } \{0..1\})$
using *test2[OF s fs uv]*
by *auto*
have $arg: \bigwedge x. (\sum_{n \in \text{basis}} (v - u) * (f (g ((v - u) * x + u)) \cdot n) * (\text{vector-derivative } g \text{ (at } ((v - u) * x + u) \text{ within } \{0..1\}) \cdot n))$
 $= (\sum_{b \in \text{basis}} f (g ((v - u) * x + u)) \cdot b * (v - u) * (\text{vector-derivative } g \text{ (at } ((v - u) * x + u) \text{ within } \{0..1\}) \cdot b))$
by (*simp add: mult.commute*)
have $((\lambda x. \sum_{b \in \text{basis}} f (g ((v - u) * x + u)) \cdot b * (\text{vector-derivative } (\lambda x. g ((v - u) * x + u)) \text{ (at } x \text{ within } \{0..1\}) \cdot b)) \text{has-integral}$
 $(\text{integral } \{u..v\} (\lambda x. \sum_{b \in \text{basis}} f (g (x)) \cdot b * (\text{vector-derivative } g \text{ (at } x \text{ within } \{0..1\}) \cdot b)))) \{0..1\}$
apply (*cut-tac Henstock-Kurzweil-Integration.has-integral-mult-right [OF *,*
where $c = v - u]$)
using *fs assms*
apply (*simp add: False subpath-def line-integral-exists-def*)
apply (*rule-tac S = (\lambda t. ((v - u) *_R t + u)) - ' s in has-integral-spike-finite*)
apply (*auto simp: inj-on-def False vd finite-vimageI scaleR-conv-of-real*
Groups-Big.sum-distrib-left
mult.assoc[symmetric] arg)
done
then show (*line-integral-exists f basis (subpath u v g)*)
by (*auto simp add: line-integral-exists-def subpath-def integrable-on-def*)
qed

type-synonym *path* = *real* \Rightarrow (*real* * *real*)
type-synonym *one-cube* = (*real* \Rightarrow (*real* * *real*))
type-synonym *one-chain* = (*int* * *path*) *set*
type-synonym *two-cube* = (*real* * *real*) \Rightarrow (*real* * *real*)
type-synonym *two-chain* = *two-cube* *set*

definition *one-chain-line-integral* :: ((real * real) \Rightarrow (real * real)) \Rightarrow ((real*real) set) \Rightarrow one-chain \Rightarrow real **where**
one-chain-line-integral F b C \equiv ($\sum (k,g) \in C. k * (\text{line-integral } F \text{ b } g)$)

definition *boundary-chain* **where**
boundary-chain s \equiv ($\forall (k, \gamma) \in s. k = 1 \vee k = -1$)

fun *coeff-cube-to-path*::(int * one-cube) \Rightarrow path
where *coeff-cube-to-path* (k, γ) = (if k = 1 then γ else (reversepath γ))

fun *rec-join* :: (int*path) list \Rightarrow path **where**
rec-join [] = ($\lambda x. 0$) |
rec-join [oneC] = *coeff-cube-to-path* oneC |
rec-join (oneC # xs) = *coeff-cube-to-path* oneC +++ (rec-join xs)

fun *valid-chain-list* **where**
valid-chain-list [] = True |
valid-chain-list [oneC] = True |
valid-chain-list (oneC # l) = (pathfinish (*coeff-cube-to-path* (oneC))) = pathstart (rec-join l) \wedge *valid-chain-list* l

lemma *joined-is-valid*:

assumes *boundary-chain*: *boundary-chain* (set l) **and**
valid-path: $\bigwedge k \gamma. (k, \gamma) \in \text{set } l \implies \text{valid-path } \gamma$ **and**
valid-chain-list-ass: *valid-chain-list* l
shows *valid-path* (rec-join l)
using *assms*
proof (*induction* l)
case Nil
then show ?case
using C1-differentiable-imp-piecewise[OF C1-differentiable-on-const[of 0 {0..1}]]
by (*auto simp add: valid-path-def*)
next
case (Cons a l)
have *: *valid-path* (rec-join ((k::int, γ) # l))
if *boundary-chain* (set (l))
 $(\bigwedge k' \gamma'. (k', \gamma') \in \text{set } l \implies \text{valid-path } \gamma')$
valid-chain-list l
valid-path (rec-join l)
 $(\bigwedge k' \gamma'. (k', \gamma') \in \text{set } ((k, \gamma) \# l) \implies \text{valid-path } \gamma')$
valid-chain-list ((k, γ) # l)
boundary-chain (set ((k, γ) # l)) **for** k γ l
proof (*cases* l = [])
case True
with that show *valid-path* (rec-join ((k, γ) # l))
by *auto*
next


```

case False
then obtain  $h\ l'$  where  $l\text{-empty}: l = h\#l'$ 
  by (meson rec-join.elims)
then show  $\text{valid-path } (\text{rec-join } ((k, \gamma) \# l))$ 
proof (simp, intro conjI impI)
  assume  $k\text{-eq-1}: k = 1$ 
  have  $0:\text{valid-path } \gamma$ 
    using that by auto
  have  $1:\text{pathfinish } \gamma = \text{pathstart } (\text{rec-join } (h\#l'))$ 
    using that(6) k-eq-1 l-empty by auto
  show  $\text{valid-path } (\gamma +++ \text{rec-join } (h\#l'))$ 
    using  $0\ 1\ \text{valid-path-join}$  that(4) l-empty by auto
next
  assume  $k \neq 1$ 
  then have  $k\text{-eq-neg-1}: k = -1$ 
    using that(7)
  by (auto simp add: boundary-chain-def)
  have  $\text{valid-path } \gamma$ 
    using that by auto
  then have  $0:\ \text{valid-path } (\text{reversepath } \gamma)$ 
    using valid-path-imp-reverse
  by auto
  have  $1:\ \text{pathfinish } (\text{reversepath } \gamma) = \text{pathstart } (\text{rec-join } (h\#l'))$ 
    using that(6) k-eq-neg-1 l-empty by auto
  show  $\text{valid-path } ((\text{reversepath } \gamma) +++ \text{rec-join } (h\#l'))$ 
    using  $0\ 1\ \text{valid-path-join}$  that(4) l-empty by blast
qed
qed
have  $\forall ps.\ \text{valid-chain-list } ps \vee (\exists i\ f\ p\ psa.\ ps = (i, f) \# p \# psa \wedge ((i = 1 \wedge \text{pathfinish } f \neq \text{pathstart } (\text{rec-join } (p \# psa))) \vee i \neq 1 \wedge \text{pathfinish } (\text{reversepath } f) \neq \text{pathstart } (\text{rec-join } (p \# psa))) \vee \neg \text{valid-chain-list } (p \# psa))$ 
  by (smt coeff-cube-to-path.elims valid-chain-list.elims(3))
moreover have  $\text{boundary-chain } (\text{set } l)$ 
  by (meson Cons.premis(1) boundary-chain-def set-subset-Cons subset-eq)
ultimately show ?case
  using * Cons by (metis (no-types) list.set-intros(2) prod.collapse valid-chain-list.simps(3))
qed

lemma pathstart-rec-join-1:
   $\text{pathstart } (\text{rec-join } ((1, \gamma) \# l)) = \text{pathstart } \gamma$ 
proof (cases l = [])
  case True
  then show  $\text{pathstart } (\text{rec-join } ((1, \gamma) \# l)) = \text{pathstart } \gamma$ 
    by simp
next
  case False
  then obtain  $h\ l'$  where  $l = h\#l'$ 
  by (meson rec-join.elims)
  then show  $\text{pathstart } (\text{rec-join } ((1, \gamma) \# l)) = \text{pathstart } \gamma$ 

```

by *simp*
qed

lemma *pathstart-rec-join-2*:

$pathstart (rec-join ((-1, \gamma) \# l)) = pathstart (reversepath \gamma)$

proof *cases*

assume $l = []$

then show $pathstart (rec-join ((-1, \gamma) \# l)) = pathstart (reversepath \gamma)$

by *simp*

next

assume $l \neq []$

then obtain $h l'$ where $l = h \# l'$

by (*meson rec-join.elims*)

then show $pathstart (rec-join ((-1, \gamma) \# l)) = pathstart (reversepath \gamma)$

by *simp*

qed

lemma *pathstart-rec-join*:

$pathstart (rec-join ((1, \gamma) \# l)) = pathstart \gamma$

$pathstart (rec-join ((-1, \gamma) \# l)) = pathstart (reversepath \gamma)$

using *pathstart-rec-join-1 pathstart-rec-join-2*

by *auto*

lemma *line-integral-exists-on-rec-join*:

assumes *boundary-chain*: *boundary-chain* (set l) and

valid-chain-list: *valid-chain-list* l and

valid-path: $\bigwedge k \gamma. (k, \gamma) \in \text{set } l \implies \text{valid-path } \gamma$ and

line-integral-exists: $\forall (k, \gamma) \in \text{set } l. \text{line-integral-exists } F \text{ basis } \gamma$

shows *line-integral-exists* $F \text{ basis } (rec-join l)$

using *assms*

proof (*induction l*)

case *Nil*

then show *?case*

proof (*simp add: line-integral-exists-def*)

have $\forall x. (\text{vector-derivative } (\lambda x. 0) \text{ (at } x)) = 0$

using *Derivative.vector-derivative-const-at*

by *auto*

then have $\forall x. ((\lambda x. 0) \text{ has-vector-derivative } 0) \text{ (at } x)$

using *Derivative.vector-derivative-const-at*

by *auto*

then have $\forall x. ((\lambda x. 0) \text{ has-vector-derivative } 0) \text{ (at } x \text{ within } \{0..1\})$

using *Derivative.vector-derivative-const-at*

by *auto*

then have $0: \forall x \in \{0..1\}. (\text{vector-derivative } (\lambda x. 0) \text{ (at } x \text{ within } \{0..1\})) = 0$

by (*simp add: gamma-deriv-at-within*)

have $(\forall x \in \{0..1\}. (\sum b \in \text{basis}. F 0 \cdot b * (\text{vector-derivative } (\lambda x. 0) \text{ (at } x \text{ within } \{0..1\}) \cdot b)) = 0)$

by (*simp add: 0*)

then have $((\lambda x. \sum b \in \text{basis}. F 0 \cdot b * (\text{vector-derivative } (\lambda x. 0) \text{ (at } x \text{ within } \{0..1\}))) = 0)$

```

{0..1}) · b)) has-integral 0) {0..1}
  by (meson has-integral-is-0)
  then show (λx. ∑ b∈basis. F 0 · b * (vector-derivative (λx. 0) (at x within
{0..1}) · b)) integrable-on {0..1}
  by auto
qed
next
case (Cons a l)
obtain k γ where aeq: a = (k,γ)
  by force
show ?case
  unfolding aeq
proof cases
  assume l-empty: l = []
  then show line-integral-exists F basis (rec-join ((k, γ) # l))
    using Cons.prem1 aeq line-integral-on-reverse-path(2) by fastforce
next
  assume l ≠ []
  then obtain h l' where l-nempty: l = h#l'
    by (meson rec-join.elims)
  show line-integral-exists F basis (rec-join ((k, γ) # l))
  proof (auto simp add: l-nempty)
    assume k-eq-1: k = 1
    have 0: line-integral-exists F basis γ
      using Cons.prem4 aeq by auto
    have 1: line-integral-exists F basis (rec-join l)
      by (metis (mono-tags) Cons boundary-chain-def list.set-intros(2) valid-chain-list.elims(3)
valid-chain-list.simps(3))
    have 2: valid-path γ
      using Cons aeq by auto
    have 3: valid-path (rec-join l)
      by (metis (no-types) Cons.prem1 boundary-chain-def joined-is-valid l-nempty
set-subset-Cons subsetCE valid-chain-list.simps(3))
    show line-integral-exists F basis (γ +++ rec-join (h#l'))
      using line-integral-distrib(2)[OF 0 1 2 3] assms l-nempty by auto
  next
    assume k ≠ 1
    then have k-eq-neg-1: k = -1
      using Cons aeq by (simp add: boundary-chain-def)
    have gamma-valid: valid-path γ
      using Cons aeq by auto
    then have 2: valid-path (reversepath γ)
      using valid-path-imp-reverse by auto
    have line-integral-exists F basis γ
      using Cons aeq by auto
    then have 0: line-integral-exists F basis (reversepath γ)
      using line-integral-on-reverse-path(2) gamma-valid
      by auto
    have 1: line-integral-exists F basis (rec-join l)

```

```

    using Cons aeq
  by (metis (mono-tags) boundary-chain-def insert-iff list.set(2) list.set-intros(2)
    valid-chain-list.elims(3) valid-chain-list.simps(3))
  have 3:valid-path (rec-join l)
    by (metis (no-types) Cons.prem(1) Cons.prem(2) Cons.prem(3) bound-
      ary-chain-def joined-is-valid l-empty set-subset-Cons subsetCE valid-chain-list.simps(3))
  show line-integral-exists F basis ((reversepath  $\gamma$ ) +++ rec-join (h#l'))
    using line-integral-distrib(2)[OF 0 1 2 3] assms l-empty
  by auto
qed
qed
qed

```

```

lemma line-integral-exists-rec-join-cons:
  assumes line-integral-exists F basis (rec-join ((1, $\gamma$ ) # l))
    ( $\bigwedge k' \gamma'. (k', \gamma') \in \text{set } ((1, \gamma) \# l) \implies \text{valid-path } \gamma'$ )
    finite basis
  shows line-integral-exists F basis ( $\gamma$  +++ (rec-join l))
proof cases
  assume l-empty: l = []
  show line-integral-exists F basis ( $\gamma$  +++ rec-join l)
    using assms(2) line-integral-distrib(2)[OF assms(1) line-integral-exists-on-degenerate-path[OF
  assms(3)], of 0]
    using degenerate-path-is-valid-path
  by (fastforce simp add: l-empty)
next
  assume l  $\neq$  []
  then obtain h l' where l = h#l'
    by (meson rec-join.elims)
  then show line-integral-exists F basis ( $\gamma$  +++ rec-join l)
    using assms by auto
qed

```

```

lemma line-integral-exists-rec-join-cons-2:
  assumes line-integral-exists F basis (rec-join ((-1, $\gamma$ ) # l))
    ( $\bigwedge k' \gamma'. (k', \gamma') \in \text{set } ((1, \gamma) \# l) \implies \text{valid-path } \gamma'$ )
    finite basis
  shows line-integral-exists F basis ((reversepath  $\gamma$ ) +++ (rec-join l))
proof cases
  assume l-empty: l = []
  show line-integral-exists F basis ((reversepath  $\gamma$ ) +++ rec-join l)
    using assms(2) line-integral-distrib(2)[OF assms(1) line-integral-exists-on-degenerate-path[OF
  assms(3)], of 0]
    using degenerate-path-is-valid-path
  by (auto simp add: l-empty)
next
  assume l  $\neq$  []
  then obtain h l' where l = h#l'
    by (meson rec-join.elims)

```

```

with assms show line-integral-exists F basis ((reversepath  $\gamma$ ) +++ rec-join l)
using assms by auto
qed

lemma line-integral-exists-on-rec-join':
assumes boundary-chain: boundary-chain (set l) and
valid-chain-list: valid-chain-list l and
valid-path:  $\bigwedge k \gamma. (k, \gamma) \in \text{set } l \implies \text{valid-path } \gamma$  and
line-integral-exists: line-integral-exists F basis (rec-join l) and
finite-basis: finite basis
shows  $\forall (k, \gamma) \in \text{set } l. \text{line-integral-exists F basis } \gamma$ 
using assms
proof (induction l)
case Nil
show ?case
by (simp add: line-integral-exists-def)
next
case ass: (Cons a l)
obtain k  $\gamma$  where k-gamma:a = (k, $\gamma$ )
by fastforce
show ?case
apply (auto simp add: k-gamma)
proof –
show line-integral-exists F basis  $\gamma$ 
proof(cases k = 1)
assume k-eq-1: k = 1
have 0: line-integral-exists F basis ( $\gamma$  +++ (rec-join l))
using line-integral-exists-rec-join-cons k-eq-1 k-gamma ass(4) ass(5) ass(6)
by auto
have 2: valid-path  $\gamma$ 
using ass k-gamma by auto
show line-integral-exists F basis  $\gamma$ 
using line-integral-exists-joinDI[OF 0 2]
by auto
next
assume k  $\neq$  1
then have k-eq-neg-1: k = -1
using ass k-gamma
by (simp add: boundary-chain-def)
have 0: line-integral-exists F basis ((reversepath  $\gamma$ ) +++ (rec-join l))
using line-integral-exists-rec-join-cons-2[OF ] k-eq-neg-1 k-gamma ass(4)
ass(5) ass(6)
by fastforce
have gamma-valid:
valid-path  $\gamma$ 
using ass k-gamma by auto
then have 2: valid-path (reversepath  $\gamma$ )
using valid-path-imp-reverse by auto
have line-integral-exists F basis (reversepath  $\gamma$ )

```

```

    using line-integral-exists-joinD1[OF 0 2] by auto
  then show line-integral-exists F basis  $(\gamma)$ 
    using line-integral-on-reverse-path(2)[OF 2] reversepath-reversepath
    by auto
qed
next
have 0: boundary-chain (set l)
  using ass(2)
  by (auto simp add: boundary-chain-def)
have 1: valid-chain-list l
  using ass(3)
  apply (auto simp add: k-gamma)
  by (metis valid-chain-list.elims(3) valid-chain-list.simps(3))
have 2:  $(\bigwedge k \gamma. (k, \gamma) \in \text{set } l \implies \text{valid-path } \gamma)$ 
  using ass(4) by auto
have 3: valid-path (rec-join l)
  using joined-is-valid[OF 0] 1 2 by auto
have 4: line-integral-exists F basis (rec-join l)
proof (cases k = 1)
  assume k-eq-1: k = 1
  have 0: line-integral-exists F basis  $(\gamma \text{ +++ } (\text{rec-join } l))$ 
    using line-integral-exists-rec-join-cons k-eq-1 k-gamma ass(4) ass(5) ass(6)
  by auto
  show line-integral-exists F basis (rec-join l)
    using line-integral-exists-joinD2[OF 0 3] by auto
next
  assume k  $\neq$  1
  then have k-eq-neg-1: k = -1
    using ass k-gamma
    by (simp add: boundary-chain-def)
  have 0: line-integral-exists F basis  $((\text{reversepath } \gamma) \text{ +++ } (\text{rec-join } l))$ 
    using line-integral-exists-rec-join-cons-2[OF ] k-eq-neg-1 k-gamma ass(4)
  ass(5) ass(6)
  by fastforce
  show line-integral-exists F basis (rec-join l)
    using line-integral-exists-joinD2[OF 0 3]
    by auto
qed
show  $\bigwedge a b. (a, b) \in \text{set } l \implies \text{line-integral-exists F basis } b$ 
  using 0 1 2 3 4 ass(1)[OF 0 1 2] ass(6)
  by fastforce
qed
qed

inductive chain-subdiv-path
  where I: chain-subdiv-path  $\gamma$  (set l) if distinct l rec-join l =  $\gamma$  valid-chain-list l

lemma valid-path-equiv-valid-chain-list:
  assumes path-eq-chain: chain-subdiv-path  $\gamma$  one-chain

```

and *boundary-chain one-chain* $\forall (k, \gamma) \in \text{one-chain. valid-path } \gamma$
shows *valid-path* γ
proof –
obtain l **where** *l-props: set* $l = \text{one-chain distinct } l \text{ rec-join } l = \gamma \text{ valid-chain-list } l$
using *chain-subdiv-path.cases path-eq-chain* **by** *force*
show *valid-path* γ
using *joined-is-valid assms l-props* **by** *blast*
qed

lemma *line-integral-rec-join-cons:*

assumes *line-integral-exists F basis* γ
line-integral-exists F basis (rec-join ((l)))
 $(\bigwedge k' \gamma'. (k', \gamma') \in \text{set } ((1, \gamma) \# l) \implies \text{valid-path } \gamma')$
finite basis
shows *line-integral F basis (rec-join ((1, \gamma) \# l)) = line-integral F basis (\gamma +++ (rec-join l))*

proof *cases*

assume *l-empty: l = []*
show *line-integral F basis (rec-join ((1, \gamma) \# l)) = line-integral F basis (\gamma +++ (rec-join l))*

using *assms line-integral-distrib(1)[OF assms(1) line-integral-exists-on-degenerate-path[OF assms(4)], of 0]*

apply *(auto simp add: l-empty)*

using *degenerate-path-is-valid-path line-integral-degenerate-path*
by *fastforce*

next

assume $l \neq []$

then obtain $h l'$ **where** *l-nempty: l = h \# l'*

by *(meson rec-join.elims)*

show *line-integral F basis (rec-join ((1, \gamma) \# l)) = line-integral F basis (\gamma +++ (rec-join l))*

using *assms* **by** *(auto simp add: l-nempty)*

qed

lemma *line-integral-rec-join-cons-2:*

assumes *line-integral-exists F basis* γ

line-integral-exists F basis (rec-join ((l)))

$(\bigwedge k' \gamma'. (k', \gamma') \in \text{set } ((-1, \gamma) \# l) \implies \text{valid-path } \gamma')$

finite basis

shows *line-integral F basis (rec-join ((-1, \gamma) \# l)) = line-integral F basis ((reversepath \gamma) +++ (rec-join l))*

proof *cases*

assume *l-empty: l = []*

have 0 : *line-integral-exists F basis (reversepath \gamma)*

using *assms line-integral-on-reverse-path(2)* **by** *fastforce*

have 1 : *valid-path (reversepath \gamma)*

using *assms* **by** *fastforce*

show *line-integral F basis (rec-join ((-1, \gamma) \# l)) = line-integral F basis ((reversepath*

```

γ) +++ (rec-join l))
  using assms line-integral-distrib(1)[OF 0 line-integral-exists-on-degenerate-path[OF
assms(4)], of 0]
  apply (auto simp add: l-empty)
  using degenerate-path-is-valid-path line-integral-degenerate-path
  by fastforce
next
  assume l ≠ []
  then obtain h l' where l-empty: l = h#l'
    by (meson rec-join.elims)
  show line-integral F basis (rec-join ((-1,γ) # l)) = line-integral F basis ((reversepath
γ) +++ (rec-join l))
    using assms by (auto simp add: l-empty)
qed

```

lemma *one-chain-line-integral-rec-join:*

```

assumes l-props: set l = one-chain distinct l valid-chain-list l and
  boundary-chain: boundary-chain one-chain and
  line-integral-exists: ∀ (k::int, γ) ∈ one-chain. line-integral-exists F basis γ and
  valid-path: ∀ (k::int, γ) ∈ one-chain. valid-path γ and
  finite-basis: finite basis
shows line-integral F basis (rec-join l) = one-chain-line-integral F basis one-chain
proof -
  have 0: sum-list (map (λ(k::int), γ). (k::int) * (line-integral F basis γ)) l) =
one-chain-line-integral F basis one-chain
    unfolding one-chain-line-integral-def
    using l-props Groups-List.comm-monoid-add-class.sum.distinct-set-conv-list[OF
l-props(2), of (λ(k, γ). (k::int) * (line-integral F basis γ))]
    by auto
  have valid-chain-list l ⇒
    boundary-chain (set l) ⇒
    (∀ (k::int, γ) ∈ set l. line-integral-exists F basis γ) ⇒
    (∀ (k::int, γ) ∈ set l. valid-path γ) ⇒
    line-integral F basis (rec-join l) = sum-list (map (λ(k::int, γ). k *
(line-integral F basis γ)) l)
  proof (induction l)
    case Nil
    show ?case
      unfolding line-integral-def boundary-chain-def
      apply (auto)
    proof
      have ∀ x. (vector-derivative (λx. 0) (at x)) = 0
        using Derivative.vector-derivative-const-at
        by auto
      then have ∀ x. ((λx. 0) has-vector-derivative 0) (at x)
        using Derivative.vector-derivative-const-at
        by auto
      then have ∀ x. ((λx. 0) has-vector-derivative 0) (at x within {0..1})
        using Derivative.vector-derivative-const-at

```



```

    by auto
  then have 0:  $\forall x \in \{0..1\}. (\text{vector-derivative } (\lambda x. 0) (\text{at } x \text{ within } \{0..1\})) = 0$ 
    by (metis (no-types) box-real(2) vector-derivative-within-cbox zero-less-one)
  have  $(\forall x \in \{0..1\}. (\sum_{b \in \text{basis}} F 0 \cdot b * (\text{vector-derivative } (\lambda x. 0) (\text{at } x \text{ within } \{0..1\}) \cdot b)) = 0)$ 
    by (simp add: 0)
  then show  $((\lambda x. \sum_{b \in \text{basis}} F 0 \cdot b * (\text{vector-derivative } (\lambda x. 0) (\text{at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } 0) \{0..1\}$ 
    by (meson has-integral-is-0)
  qed
next
case ass: (Cons a l)
obtain  $k::\text{int}$  and  $\gamma::\text{one-cube}$  where props:  $a = (k, \gamma)$ 
proof
  let  $?k2 = \text{fst } a$ 
  let  $?\gamma2 = \text{snd } a$ 
  show  $a = (?k2, ?\gamma2)$ 
  by auto
qed
have line-integral-exists F basis (rec-join (a # l))
  using line-integral-exists-on-rec-join[OF ass(3) ass(2)] ass(5) ass(4)
  by blast
have boundary-chain (set l)
  by (meson ass.prem(2) boundary-chain-def list.set-intros(2))
have val-l:  $\bigwedge f i. (i, f) \in \text{set } l \implies \text{valid-path } f$ 
  using ass.prem(4) by fastforce
have vcl-l: valid-chain-list l
  by (metis (no-types) ass.prem(1) valid-chain-list.elims(3) valid-chain-list.simps(3))
have line-integral-exists-on-joined:
  line-integral-exists F basis (rec-join l)
  by (metis  $\langle \text{boundary-chain } (\text{set } l) \rangle \langle \text{line-integral-exists } F \text{ basis } (\text{rec-join } (a \# l)) \rangle$ 
  emptyE val-l vcl-l joined-is-valid line-integral-exists-joinD2 line-integral-exists-on-rec-join
  list.set(1) neq-Nil-conv rec-join.simps(3))
have valid-path (rec-join (a # l))
  using joined-is-valid ass(5) ass(3) ass(2) by blast
then have joined-is-valid: valid-path (rec-join l)
  using  $\langle \text{boundary-chain } (\text{set } l) \rangle$  val-l vcl-l joined-is-valid by blast
show ?case
proof (clarsimp, cases)
  assume k-eq-1:  $(k::\text{int}) = 1$ 
  have line-integral-exists-on-gamma: line-integral-exists F basis  $\gamma$ 
    using ass props by auto
  have gamma-is-valid: valid-path  $\gamma$ 
    using ass props by auto
  have line-int-rw: line-integral F basis (rec-join  $((k, \gamma) \# l)$ ) = line-integral F
  basis  $(\gamma \text{ +++ rec-join } l)$ 
  proof -
    have gam-int: line-integral-exists F basis  $\gamma$  using ass props by auto

```

```

have rec-join-int: line-integral-exists F basis (rec-join l)
  using line-integral-exists-on-rec-join
  using line-integral-exists-on-joined by blast
show ?thesis
  using line-integral-rec-join-cons[OF gam-int rec-join-int] ass k-eq-1 fi-
nite-basis props by force
qed
show line-integral F basis (rec-join (a # l)) =
  (case a of (x, γ) ⇒ real-of-int x * line-integral F basis γ) + (∑ (x, γ)←l.
real-of-int x * line-integral F basis γ)
  apply (simp add: props line-int-rw)
  using line-integral-distrib[OF line-integral-exists-on-gamma line-integral-exists-on-joined
gamma-is-valid joined-is-valid]
  ass k-eq-1 vcl-l
  by (auto simp: boundary-chain-def props)
next
assume k ≠ 1
then have k-eq-neg-1: k = -1
  using ass props
  by (auto simp add: boundary-chain-def)
have line-integral-exists-on-gamma:
  line-integral-exists F basis (reversepath γ)
  using line-integral-on-reverse-path ass props
  by auto
have gamma-is-valid: valid-path (reversepath γ)
  using valid-path-imp-reverse ass props by auto
have line-int-rw: line-integral F basis (rec-join ((k, γ) # l)) = line-integral F
basis ((reversepath γ) +++ rec-join l)
proof -
  have gam-int: line-integral-exists F basis γ using ass props by auto
  have rec-join-int: line-integral-exists F basis (rec-join l)
  using line-integral-exists-on-rec-join
  using line-integral-exists-on-joined by blast
  show ?thesis
  using line-integral-rec-join-cons-2[OF gam-int rec-join-int]
  using ass k-eq-neg-1
  using finite-basis props by blast
qed
show line-integral F basis (rec-join (a # l)) =
  (case a of (x, γ) ⇒ real-of-int x * line-integral F basis γ) + (∑ (x, γ)←l.
real-of-int x * line-integral F basis γ)
  apply (simp add: props line-int-rw)
  using line-integral-distrib[OF line-integral-exists-on-gamma line-integral-exists-on-joined
gamma-is-valid joined-is-valid]
  props ass line-integral-on-reverse-path(1)[of γ F basis] k-eq-neg-1
  using ⟨boundary-chain (set l)⟩ vcl-l by auto
qed
qed
then have 1: line-integral F basis (rec-join l) = sum-list (map (λ(k::int, γ). k *

```

(*line-integral F basis γ*) l)
using *l-props assms* **by** *auto*
then show *?thesis*
using *0 1* **by** *auto*
qed

lemma *line-integral-on-path-eq-line-integral-on-equiv-chain*:
assumes *path-eq-chain: chain-subdiv-path γ one-chain* **and**
boundary-chain: boundary-chain one-chain **and**
line-integral-exists: $\forall (k::int, \gamma) \in one-chain. line-integral-exists F basis \gamma$ **and**
valid-path: $\forall (k::int, \gamma) \in one-chain. valid-path \gamma$ **and**
finite-basis: finite basis
shows *one-chain-line-integral F basis one-chain = line-integral F basis γ*
line-integral-exists F basis γ
valid-path γ

proof –
obtain *l* **where** *l-props: set l = one-chain distinct l rec-join l = γ valid-chain-list*
l
using *chain-subdiv-path.cases path-eq-chain* **by** *force*
show *line-integral-exists F basis γ*
using *line-integral-exists-on-rec-join assms l-props*
by *blast*
show *valid-path γ*
using *joined-is-valid assms l-props*
by *blast*
have *line-integral F basis (rec-join l) = one-chain-line-integral F basis one-chain*
using *one-chain-line-integral-rec-join l-props assms* **by** *auto*
then show *one-chain-line-integral F basis one-chain = line-integral F basis γ*
using *l-props*
by *auto*
qed

lemma *line-integral-on-path-eq-line-integral-on-equiv-chain'*:
assumes *path-eq-chain: chain-subdiv-path γ one-chain* **and**
boundary-chain: boundary-chain one-chain **and**
line-integral-exists: line-integral-exists F basis γ **and**
valid-path: $\forall (k, \gamma) \in one-chain. valid-path \gamma$ **and**
finite-basis: finite basis
shows *one-chain-line-integral F basis one-chain = line-integral F basis γ*
 $\forall (k, \gamma) \in one-chain. line-integral-exists F basis \gamma$
proof –
obtain *l* **where** *l-props: set l = one-chain distinct l rec-join l = γ valid-chain-list*
l
using *chain-subdiv-path.cases path-eq-chain* **by** *force*
show *0: $\forall (k, \gamma) \in one-chain. line-integral-exists F basis \gamma$*
using *line-integral-exists-on-rec-join' assms l-props*
by *blast*
show *one-chain-line-integral F basis one-chain = line-integral F basis γ*
using *line-integral-on-path-eq-line-integral-on-equiv-chain(1)[OF assms(1) assms(2)]*

0 *assms*(4) *assms*(5)] **by auto**
qed

definition *chain-subdiv-chain* **where**

chain-subdiv-chain one-chain1 subdiv
 $\equiv \exists f. (\bigcup (f \text{ ' one-chain1})) = \text{subdiv} \wedge$
 $(\forall c \in \text{one-chain1}. \text{chain-subdiv-path } (\text{coeff-cube-to-path } c) (f \ c)) \wedge$
 $\text{pairwise } (\lambda p \ p'. f \ p \cap f \ p' = \{\}) \text{ one-chain1} \wedge$
 $(\forall x \in \text{one-chain1}. \text{finite } (f \ x))$

lemma *chain-subdiv-chain-character*:

shows *chain-subdiv-chain one-chain1 subdiv* \longleftrightarrow
 $(\exists f. \bigcup (f \text{ ' one-chain1}) = \text{subdiv} \wedge$
 $(\forall (k, \gamma) \in \text{one-chain1}.$
 $\text{if } k = 1$
 $\text{then } \text{chain-subdiv-path } \gamma (f (k, \gamma))$
 $\text{else } \text{chain-subdiv-path } (\text{reversepath } \gamma) (f (k, \gamma))) \wedge$
 $(\forall p \in \text{one-chain1}.$
 $\forall p' \in \text{one-chain1}. p \neq p' \longrightarrow f \ p \cap f \ p' = \{\}) \wedge$
 $(\forall x \in \text{one-chain1}. \text{finite } (f \ x)))$

unfolding *chain-subdiv-chain-def*

by (*safe*; *intro exI conjI iffI*; *fastforce simp add: pairwise-def*)

lemma *chain-subdiv-chain-imp-finite-subdiv*:

assumes *finite one-chain1*
chain-subdiv-chain one-chain1 subdiv
shows *finite subdiv*
using *assms* **by** (*auto simp add: chain-subdiv-chain-def*)

lemma *valid-subdiv-imp-valid-one-chain*:

assumes *chain1-eq-chain2*: *chain-subdiv-chain one-chain1 subdiv* **and**
boundary-chain1: *boundary-chain one-chain1* **and**
boundary-chain2: *boundary-chain subdiv* **and**
valid-path: $\forall (k, \gamma) \in \text{subdiv}. \text{valid-path } \gamma$
shows $\forall (k, \gamma) \in \text{one-chain1}. \text{valid-path } \gamma$

proof –

obtain *f* **where** *f*-*props*:

$((\bigcup (f \text{ ' one-chain1})) = \text{subdiv})$
 $(\forall (k, \gamma) \in \text{one-chain1}. \text{if } k = 1 \text{ then } \text{chain-subdiv-path } \gamma (f (k, \gamma)) \text{ else } \text{chain-subdiv-path}$
 $(\text{reversepath } \gamma) (f (k, \gamma)))$
 $(\forall p \in \text{one-chain1}. \forall p' \in \text{one-chain1}. p \neq p' \longrightarrow f \ p \cap f \ p' = \{\})$

using *chain1-eq-chain2 chain-subdiv-chain-character* **by auto**

have $\bigwedge k \ \gamma. (k, \gamma) \in \text{one-chain1} \implies \text{valid-path } \gamma$

proof –

fix *k* γ

assume *ass*: $(k, \gamma) \in \text{one-chain1}$

show *valid-path* γ

proof (*cases* $k = 1$)

assume *k-eq-1*: $k = 1$

```

then have i:chain-subdiv-path  $\gamma$  ( $f(k,\gamma)$ )
  using f-props(2) ass by auto
have ii:boundary-chain ( $f(k,\gamma)$ )
  using f-props(1) ass assms
  apply (simp add: boundary-chain-def)
  by blast
have iv:  $\forall (k, \gamma) \in f (k, \gamma). \text{valid-path } \gamma$ 
  using f-props(1) ass assms
  by blast
show ?thesis
  using valid-path-equiv-valid-chain-list[OF i ii iv]
  by auto
next
assume  $k \neq 1$ 
then have k-eq-neg1:  $k = -1$ 
  using ass boundary-chain1
  by (auto simp add: boundary-chain-def)
then have i:chain-subdiv-path (reversepath  $\gamma$ ) ( $f(k,\gamma)$ )
  using f-props(2) ass using  $\langle k \neq 1 \rangle$  by auto
have ii:boundary-chain ( $f(k,\gamma)$ )
  using f-props(1) ass assms by (auto simp add: boundary-chain-def)
have iv:  $\forall (k, \gamma) \in f (k, \gamma). \text{valid-path } \gamma$ 
  using f-props(1) ass assms
  by blast
have valid-path (reversepath  $\gamma$ )
  using valid-path-equiv-valid-chain-list[OF i ii iv]
  by auto
then show ?thesis
  using reversepath-reversepath valid-path-imp-reverse
  by force
qed
qed
then show valid-path1:  $\forall (k, \gamma) \in \text{one-chain1}. \text{valid-path } \gamma$ 
  by auto
qed

lemma one-chain-line-integral-eq-line-integral-on-sudivision:
  assumes chain1-eq-chain2: chain-subdiv-chain one-chain1 subdiv and
  boundary-chain1: boundary-chain one-chain1 and
  boundary-chain2: boundary-chain subdiv and
  line-integral-exists-on-chain2:  $\forall (k, \gamma) \in \text{subdiv}. \text{line-integral-exists } F \text{ basis } \gamma$ 
and
  valid-path:  $\forall (k, \gamma) \in \text{subdiv}. \text{valid-path } \gamma$  and
  finite-chain1: finite one-chain1 and
  finite-basis: finite basis
shows one-chain-line-integral F basis one-chain1 = one-chain-line-integral F
basis subdiv
   $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$ 
proof –

```

```

obtain  $f$  where  $f$ -props:
  (( $\bigcup (f \text{ ' one-chain1})$ ) = subdiv)
  ( $\forall (k,\gamma) \in \text{one-chain1}$ . if  $k = 1$  then chain-subdiv-path  $\gamma (f(k,\gamma))$  else chain-subdiv-path
(reversepath  $\gamma$ ) ( $f(k,\gamma)$ ))
  ( $\forall p \in \text{one-chain1}$ .  $\forall p' \in \text{one-chain1}$ .  $p \neq p' \longrightarrow f p \cap f p' = \{\}$ )
  ( $\forall x \in \text{one-chain1}$ . finite ( $f x$ ))
  using chain1-eq-chain2 chain-subdiv-chain-character by (auto simp add: pair-
wise-def chain-subdiv-chain-def)
  then have 0: one-chain-line-integral  $F$  basis subdiv = one-chain-line-integral  $F$ 
basis ( $\bigcup (f \text{ ' one-chain1})$ )
  by auto
have finite-chain2: finite subdiv
  using finite-chain1  $f$ -props(1)  $f$ -props(4)
  apply (simp add: image-def)
  using  $f$ -props(1) by auto
have  $\bigwedge k \gamma$ .  $(k,\gamma) \in \text{one-chain1} \implies \text{line-integral-exists } F \text{ basis } \gamma$ 
proof -
  fix  $k \gamma$ 
  assume  $ass$ :  $(k,\gamma) \in \text{one-chain1}$ 
  show line-integral-exists  $F$  basis  $\gamma$ 
  proof (cases  $k = 1$ )
    assume  $k$ -eq-1:  $k = 1$ 
    then have  $i$ : chain-subdiv-path  $\gamma (f(k,\gamma))$ 
      using  $f$ -props(2)  $ass$  by auto
    have  $ii$ : boundary-chain ( $f(k,\gamma)$ )
      using  $f$ -props(1)  $ass$   $assms$  by (auto simp add: boundary-chain-def)
    have  $iii$ :  $\forall (k, \gamma) \in f (k, \gamma)$ . line-integral-exists  $F$  basis  $\gamma$ 
      using  $f$ -props(1)  $ass$   $assms$ 
      by blast
    have  $iv$ :  $\forall (k, \gamma) \in f (k, \gamma)$ . valid-path  $\gamma$ 
      using  $f$ -props(1)  $ass$   $assms$ 
      by blast
    show ?thesis
      using line-integral-on-path-eq-line-integral-on-equiv-chain(2)[OF  $i$   $ii$   $iii$   $iv$ 
finite-basis]
      by auto
  next
  assume  $k \neq 1$ 
  then have  $k$ -eq-neg1:  $k = -1$ 
  using  $ass$  boundary-chain1
  by (auto simp add: boundary-chain-def)
  then have  $i$ : chain-subdiv-path (reversepath  $\gamma$ ) ( $f(k,\gamma)$ )
  using  $f$ -props(2)  $ass$  by auto
  have  $ii$ : boundary-chain ( $f(k,\gamma)$ )
  using  $f$ -props(1)  $ass$   $assms$  by (auto simp add: boundary-chain-def)
  have  $iii$ :  $\forall (k, \gamma) \in f (k, \gamma)$ . line-integral-exists  $F$  basis  $\gamma$ 
  using  $f$ -props(1)  $ass$   $assms$ 
  by blast
  have  $iv$ :  $\forall (k, \gamma) \in f (k, \gamma)$ . valid-path  $\gamma$ 

```

```

    using f-props(1) ass assms
    by blast
  have x: line-integral-exists F basis (reversepath  $\gamma$ )
    using line-integral-on-path-eq-line-integral-on-equiv-chain(2)[OF i ii iii iv
finite-basis]
    by auto
  have valid-path (reversepath  $\gamma$ )
    using line-integral-on-path-eq-line-integral-on-equiv-chain(3)[OF i ii iii iv
finite-basis]
    by auto
  then show ?thesis
    using line-integral-on-reverse-path(2) reversepath-reversepath x
    by fastforce
qed
qed
then show line-integral-exists-on-chain1:  $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists}$ 
F basis  $\gamma$ 
  by auto
have 1: one-chain-line-integral F basis ( $\bigcup (f \text{ ' one-chain1})$ ) = one-chain-line-integral
F basis one-chain1
proof -
  have 0: one-chain-line-integral F basis ( $\bigcup (f \text{ ' one-chain1})$ ) =
    ( $\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral F}$ 
basis one-chain)
  proof -
    have finite:  $\forall \text{chain} \in (f \text{ ' one-chain1}). \text{finite chain}$ 
    using f-props(1) finite-chain2
    by (meson Sup-upper finite-subset)
    have disj:  $\forall A \in f \text{ ' one-chain1}. \forall B \in f \text{ ' one-chain1}. A \neq B \longrightarrow A \cap B = \{\}$ 
    by (metis (no-types, opaque-lifting) f-props(3) image-iff)
    show one-chain-line-integral F basis ( $\bigcup (f \text{ ' one-chain1})$ ) =
      ( $\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral}$ 
F basis one-chain)
    using Groups-Big.comm-monoid-add-class.sum.Union-disjoint[OF finite disj]
    one-chain-line-integral-def
    by auto
  qed
  have 1: ( $\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral F basis one-chain}$ )
=
  one-chain-line-integral F basis one-chain1
  proof -
    have ( $\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral F basis one-chain}$ )
=
      ( $\sum (k, \gamma) \in \text{one-chain1}. k * (\text{line-integral F basis } \gamma)$ )
    proof -
      have i: ( $\sum \text{one-chain} \in (f \text{ ' (one-chain1 - \{p. fp = \{\})\})}. \text{one-chain-line-integral}$ 
F basis one-chain) =
      ( $\sum (k, \gamma) \in \text{one-chain1} - \{p. fp = \{\}. k * (\text{line-integral}$ 
F basis  $\gamma)$ )

```

```

proof –
  have inj-on f (one-chain1 – {p. f p = {}})
    unfolding inj-on-def using f-props(3) by blast
    then have 0: ( $\sum$  one-chain  $\in$  (f ‘ (one-chain1 – {p. f p = {}})}).
one-chain-line-integral F basis one-chain)
      = ( $\sum$  (k,  $\gamma$ )  $\in$  (one-chain1 – {p. f p = {}})
      {{}}). one-chain-line-integral F basis (f (k,  $\gamma$ )))
    using Groups-Big.comm-monoid-add-class.sum.reindex
    by auto
  have  $\bigwedge$  k  $\gamma$ . (k,  $\gamma$ )  $\in$  (one-chain1 – {p. f p = {}})  $\implies$ 
    one-chain-line-integral F basis (f(k,  $\gamma$ )) = k * (line-integral F
basis  $\gamma$ )
proof –
  fix k  $\gamma$ 
  assume ass: (k,  $\gamma$ )  $\in$  (one-chain1 – {p. f p = {}})
  have bchain: boundary-chain (f(k, $\gamma$ ))
    using f-props(1) boundary-chain2 ass
    by (auto simp add: boundary-chain-def)
  have wexist:  $\forall$  (k,  $\gamma$ ) $\in$ (f(k, $\gamma$ )). line-integral-exists F basis  $\gamma$ 
    using f-props(1) line-integral-exists-on-chain2 ass
    by blast
  have vpath:  $\forall$  (k,  $\gamma$ ) $\in$ (f(k,  $\gamma$ )). valid-path  $\gamma$ 
    using f-props(1) assms ass
    by blast
  show one-chain-line-integral F basis (f (k,  $\gamma$ )) = k * line-integral F basis
 $\gamma$ 
proof(cases k = 1)
  assume k-eq-1: k = 1
  have one-chain-line-integral F basis (f (k,  $\gamma$ )) = line-integral F basis  $\gamma$ 
  using f-props(2) k-eq-1 line-integral-on-path-eq-line-integral-on-equiv-chain
bchain wexist vpath ass finite-basis
    by auto
  then show one-chain-line-integral F basis (f (k,  $\gamma$ )) = k * line-integral
F basis  $\gamma$ 
    using k-eq-1 by auto
  next
  assume k  $\neq$  1
  then have k-eq-neg1: k = -1
    using ass boundary-chain1
    by (auto simp add: boundary-chain-def)
  have one-chain-line-integral F basis (f (k,  $\gamma$ )) = line-integral F basis
(reversepath  $\gamma$ )
    using f-props(2) k-eq-neg1 line-integral-on-path-eq-line-integral-on-equiv-chain
bchain wexist vpath ass finite-basis
    by auto
  then have one-chain-line-integral F basis (f (k,  $\gamma$ )) = - (line-integral
F basis  $\gamma$ )
    using line-integral-on-reverse-path(1) ass line-integral-exists-on-chain1
valid-subdiv-imp-valid-one-chain[OF chain1-eq-chain2 boundary-chain1

```



```

boundary-chain2 valid-path]
  by force
  then show one-chain-line-integral F basis (f (k, γ)) = k * line-integral
F basis γ
    using k-eq-neg1 by auto
  qed
  then have (∑ (k, γ) ∈ (one-chain1 - {p. f p = {}}). one-chain-line-integral
F basis (f (k, γ)))
    = (∑ (k, γ) ∈ (one-chain1 - {p. f p = {}}). k * (line-integral F
basis γ))
    by (auto intro!: Finite-Cartesian-Product.sum-cong-aux)
  then show (∑ one-chain ∈ (f ' (one-chain1 - {p. f p = {}})). one-chain-line-integral
F basis one-chain) =
    (∑ (k, γ) ∈ (one-chain1 - {p. f p =
{}}). k * (line-integral F basis γ))
    using 0 by auto
  qed
  have ∧ (k::int) γ. (k, γ) ∈ one-chain1 ⇒ (f (k, γ) = { }) ⇒ (k, γ) ∈
{(k', γ'). k' * (line-integral F basis γ') = 0}
  proof-
    fix k::int
    fix γ::one-cube
    assume ass:(k, γ) ∈ one-chain1
    (f (k, γ) = { })
    then have zero-line-integral:one-chain-line-integral F basis (f (k, γ)) = 0
    using one-chain-line-integral-def
    by auto
    have bchain: boundary-chain (f(k,γ))
    using f-props(1) boundary-chain2 ass
    by (auto simp add: boundary-chain-def)
    have wexist: ∀ (k, γ)∈(f(k,γ)). line-integral-exists F basis γ
    using f-props(1) line-integral-exists-on-chain2 ass
    by blast
    have vpath: ∀ (k, γ)∈(f(k, γ)). valid-path γ
    using f-props(1) assms ass by blast
    have one-chain-line-integral F basis (f (k, γ)) = k * line-integral F basis γ
    proof(cases k = 1)
      assume k-eq-1: k = 1
      have one-chain-line-integral F basis (f (k, γ)) = line-integral F basis γ
      using f-props(2) k-eq-1 line-integral-on-path-eq-line-integral-on-equiv-chain
bchain wexist vpath ass finite-basis
      by auto
    then show one-chain-line-integral F basis (f (k, γ)) = k * line-integral
F basis γ
      using k-eq-1
      by auto
    next
      assume k ≠ 1

```

```

then have k-eq-neg1:  $k = -1$ 
  using ass boundary-chain1
  by (auto simp add: boundary-chain-def)
  have one-chain-line-integral F basis ( $f(k, \gamma) = \text{line-integral } F \text{ basis}$ 
(reversepath  $\gamma$ ))
    using f-props(2) k-eq-neg1 line-integral-on-path-eq-line-integral-on-equiv-chain
bchain wexist vpath ass finite-basis
    by auto
  then have one-chain-line-integral F basis ( $f(k, \gamma) = -(\text{line-integral } F$ 
basis  $\gamma$ ))
    using line-integral-on-reverse-path(1) ass line-integral-exists-on-chain1
valid-subdiv-imp-valid-one-chain[OF chain1-eq-chain2 boundary-chain1
boundary-chain2 valid-path]
    by force
  then show one-chain-line-integral F basis ( $f(k::int, \gamma) = k * \text{line-integral}$ 
F basis  $\gamma$ )
    using k-eq-neg1 by auto
  qed
  then show  $(k, \gamma) \in \{(k'::int, \gamma'). k' * \text{line-integral } F \text{ basis } \gamma' = 0\}$ 
    using zero-line-integral by auto
  qed
  then have ii: $(\sum \text{one-chain} \in (f'(one-chain1 - \{p.f p = \{\}\})). \text{one-chain-line-integral}$ 
F basis one-chain}) =

$$(\sum \text{one-chain} \in (f'(one-chain1))).$$

one-chain-line-integral F basis one-chain)
    proof –
      have  $\bigwedge \text{one-chain}. \text{one-chain} \in (f'(one-chain1)) - (f'(one-chain1 - \{p.f p = \{\}\})) \implies$ 

$$\text{one-chain-line-integral } F \text{ basis}$$

one-chain} = 0
      proof –
        fix one-chain
        assume  $\text{one-chain} \in (f'(one-chain1)) - (f'(one-chain1 - \{p.f p = \{\}\}))$ 

$$\{\}\})$$

        then show one-chain-line-integral F basis one-chain} = 0
          by (auto simp add: one-chain-line-integral-def)
        qed
      then have  $0: (\sum \text{one-chain} \in f'(one-chain1) - (f'(one-chain1 - \{p.f p = \{\}\})). \text{one-chain-line-integral } F \text{ basis one-chain})$ 

$$= 0$$

        using comm-monoid-add-class.sum.neutral by auto
      then have  $(\sum \text{one-chain} \in f'(one-chain1)). \text{one-chain-line-integral } F \text{ basis one-chain}$ 

$$- (\sum \text{one-chain} \in (f'(one-chain1 - \{p.f p = \{\}\})). \text{one-chain-line-integral } F \text{ basis one-chain})$$


$$= 0$$

      proof –
        have finte: finite ( $f'(one-chain1)$ ) using finite-chain1 by auto
        show ?thesis

```

```

using Groups-Big.sum-diff[OF finite, of (f ' (one-chain1 - {p. f p =
{}}))
  one-chain-line-integral F basis]
  0
by auto
qed
then show ( $\sum$  one-chain  $\in$  (f ' (one-chain1 - {p. f p = {}})). one-chain-line-integral
F basis one-chain) =
  ( $\sum$  one-chain  $\in$  (f ' (one-chain1)).
one-chain-line-integral F basis one-chain)
by auto
qed
have  $\bigwedge$  (k::int)  $\gamma$ . (k,  $\gamma$ )  $\in$  one-chain1  $\implies$  (f (k,  $\gamma$ ) = { })  $\implies$  f (k,  $\gamma$ )  $\in$ 
{chain. one-chain-line-integral F basis chain = 0}
proof-
  fix k::int
  fix  $\gamma$ ::one-cube
  assume ass:(k,  $\gamma$ )  $\in$  one-chain1 (f (k,  $\gamma$ ) = { })
  then have one-chain-line-integral F basis (f (k,  $\gamma$ )) = 0
    using one-chain-line-integral-def by auto
  then show f (k,  $\gamma$ )  $\in$  {p'. (one-chain-line-integral F basis p' = 0)}
    by auto
qed
then have iii:( $\sum$  (k::int, $\gamma$ ) $\in$ one-chain1 - {p. f p = {}}. k*(line-integral F
basis  $\gamma$ ))
  = ( $\sum$  (k::int, $\gamma$ ) $\in$ one-chain1. k*(line-integral
F basis  $\gamma$ ))
proof-
  have  $\bigwedge$  k  $\gamma$ . (k, $\gamma$ ) $\in$ one-chain1 - (one-chain1 - {p. f p = {}})
     $\implies$  k* (line-integral F basis  $\gamma$ ) = 0
proof-
  fix k  $\gamma$ 
  assume ass: (k, $\gamma$ ) $\in$ one-chain1 - (one-chain1 - {p. f p = {}})
  then have f(k,  $\gamma$ ) = { }
    by auto
  then have one-chain-line-integral F basis (f(k,  $\gamma$ )) = 0
    by (auto simp add: one-chain-line-integral-def)
  then have zero-line-integral:one-chain-line-integral F basis (f (k,  $\gamma$ )) =
0
    using one-chain-line-integral-def by auto
  have bchain: boundary-chain (f(k, $\gamma$ ))
    using f-props(1) boundary-chain2 ass
    by (auto simp add: boundary-chain-def)
  have wexist:  $\forall$  (k,  $\gamma$ ) $\in$ (f(k, $\gamma$ )). line-integral-exists F basis  $\gamma$ 
    using f-props(1) line-integral-exists-on-chain2 ass
    by blast
  have vpath:  $\forall$  (k,  $\gamma$ ) $\in$ (f(k,  $\gamma$ )). valid-path  $\gamma$ 
    using f-props(1) assms ass
    by blast

```

```

have one-chain-line-integral F basis (f (k, γ)) = k * line-integral F basis
γ
proof(cases k = 1)
  assume k-eq-1: k = 1
  have one-chain-line-integral F basis (f (k, γ)) = line-integral F basis γ
  using f-props(2) k-eq-1 line-integral-on-path-eq-line-integral-on-equiv-chain
bchain wexist vpath ass finite-basis
  by auto
  then show one-chain-line-integral F basis (f (k, γ)) = k * line-integral
F basis γ
    using k-eq-1
    by auto
  next
  assume k ≠ 1
  then have k-eq-neg1: k = -1
  using ass boundary-chain1
  by (auto simp add: boundary-chain-def)
  have one-chain-line-integral F basis (f (k, γ)) = line-integral F basis
(reversepath γ)
  using f-props(2) k-eq-neg1 line-integral-on-path-eq-line-integral-on-equiv-chain
bchain wexist vpath ass finite-basis
  by auto
  then have one-chain-line-integral F basis (f (k, γ)) = - (line-integral
F basis γ)
  using line-integral-on-reverse-path(1) ass line-integral-exists-on-chain1
valid-subdiv-imp-valid-one-chain[OF chain1-eq-chain2 boundary-chain1
boundary-chain2 valid-path]
  by force
  then show one-chain-line-integral F basis (f (k, γ)) = k * line-integral
F basis γ
    using k-eq-neg1
    by auto
  qed
  then show k* (line-integral F basis γ) = 0
  using zero-line-integral
  by auto
  qed
  then have  $\forall (k::int, \gamma) \in \text{one-chain1} - (\text{one-chain1} - \{p. f p = \{\}\})$ .
 $k * (\text{line-integral F basis } \gamma) = 0$  by auto
  then have  $(\sum (k::int, \gamma) \in \text{one-chain1} - (\text{one-chain1} - \{p. f p = \{\}\})$ .
 $k * (\text{line-integral F basis } \gamma)) = 0$ 
  using Groups-Big.comm-monoid-add-class.sum.neutral
[of one-chain1 - (one-chain1 - {p. f p = {}}) (λ(k::int, γ). k*
(line-integral F basis γ))]
  by (simp add: split-beta)
  then have  $(\sum (k::int, \gamma) \in \text{one-chain1}$ .  $k * (\text{line-integral F basis } \gamma)) -$ 
 $(\sum (k::int, \gamma) \in (\text{one-chain1} - \{p. f p = \{\}\})$ .  $k * (\text{line-integral F}$ 
basis γ)) = 0
  using Groups-Big.sum-diff[OF finite-chain1]

```

by (metis (no-types) Diff-subset $\langle (\sum (k, \gamma) \in \text{one-chain1} - (\text{one-chain1} - \{p, f p = \{\}\})) . k * \text{line-integral } F \text{ basis } \gamma) = 0 \rangle \langle \bigwedge f B. B \subseteq \text{one-chain1} \implies \text{sum } f (\text{one-chain1} - B) = \text{sum } f \text{ one-chain1} - \text{sum } f B \rangle$)
 then show ?thesis by auto
 qed
 show ?thesis using i ii iii by auto
 qed
 then show ?thesis using one-chain-line-integral-def by auto
 qed
 show ?thesis using 0 1 by auto
 qed
 show one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis subdiv using 0 1 by auto
 qed

lemma one-chain-line-integral-eq-line-integral-on-sudivision':

assumes chain1-eq-chain2: chain-subdiv-chain one-chain1 subdiv and
 boundary-chain1: boundary-chain one-chain1 and
 boundary-chain2: boundary-chain subdiv and
 line-integral-exists-on-chain1: $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$ and
 valid-path: $\forall (k, \gamma) \in \text{subdiv}. \text{valid-path } \gamma$ and
 finite-chain1: finite one-chain1 and
 finite-basis: finite basis
 shows one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis subdiv
 $\forall (k, \gamma) \in \text{subdiv}. \text{line-integral-exists } F \text{ basis } \gamma$
proof –
 obtain f where f-props:
 (($\bigcup (f \text{ ' one-chain1})) = \text{subdiv}$)
 ($\forall (k, \gamma) \in \text{one-chain1}. \text{if } k = 1 \text{ then chain-subdiv-path } \gamma (f(k, \gamma)) \text{ else chain-subdiv-path (reversepath } \gamma) (f(k, \gamma))$)
 ($\forall p \in \text{one-chain1}. \forall p' \in \text{one-chain1}. p \neq p' \implies f p \cap f p' = \{\}$)
 ($\forall x \in \text{one-chain1}. \text{finite } (f x)$)
 using chain1-eq-chain2 chain-subdiv-chain-character by (auto simp add: pairwise-def chain-subdiv-chain-def)
 have finite-chain2: finite subdiv
 using finite-chain1 f-props(1) f-props(4) by blast
 have $\bigwedge k \gamma. (k, \gamma) \in \text{subdiv} \implies \text{line-integral-exists } F \text{ basis } \gamma$
proof –
 fix k γ
 assume ass: $(k, \gamma) \in \text{subdiv}$
 then obtain $k' \gamma'$ where kp-gammap: $(k', \gamma') \in \text{one-chain1} (k, \gamma) \in f(k', \gamma')$
 using f-props(1) by fastforce
 show line-integral-exists F basis γ
proof (cases $k' = 1$)
 assume k-eq-1: $k' = 1$
 then have i: chain-subdiv-path $\gamma' (f(k', \gamma'))$
 using f-props(2) kp-gammap ass by auto

```

have ii:boundary-chain (f(k',γ'))
using f-props(1) ass assms kp-gammap by (meson UN-I boundary-chain-def)
have iii:line-integral-exists F basis γ'
  using assms kp-gammap by blast
have iv:  $\forall (k, \gamma) \in f(k', \gamma'). \text{ valid-path } \gamma$ 
  using f-props(1) ass assms kp-gammap by blast
show ?thesis
  using line-integral-on-path-eq-line-integral-on-equiv-chain'(2)[OF i ii iii iv
finite-basis] kp-gammap
  by auto
next
assume k'  $\neq$  1
then have k-eq-neg1: k' = -1
  using boundary-chain1 kp-gammap
  by (auto simp add: boundary-chain-def)
then have i:chain-subdiv-path (reversepath γ') (f(k',γ'))
  using f-props(2) kp-gammap by auto
have ii:boundary-chain (f(k',γ'))
  using f-props(1) assms kp-gammap by (meson UN-I boundary-chain-def)
have iii:  $\forall (k, \gamma) \in f(k', \gamma'). \text{ valid-path } \gamma$ 
  using f-props(1) ass assms kp-gammap by blast
have iv: valid-path (reversepath γ')
  using valid-path-equiv-valid-chain-list[OF i ii iii]
  by force
have line-integral-exists F basis γ'
  using assms kp-gammap by blast
then have x: line-integral-exists F basis (reversepath γ')
  using iv line-integral-on-reverse-path(2) valid-path-reversepath
  by fastforce
show ?thesis
  using line-integral-on-path-eq-line-integral-on-equiv-chain'(2)[OF i ii x iii
finite-basis] kp-gammap
  by auto
qed
qed
then show  $\forall (k, \gamma) \in \text{subdiv. line-integral-exists } F \text{ basis } \gamma$  by auto
then show one-chain-line-integral F basis one-chain1 = one-chain-line-integral
F basis subdiv
  using one-chain-line-integral-eq-line-integral-on-sudivision(1) assms
  by auto
qed

lemma line-integral-sum-gen:
assumes finite-basis:
  finite basis and
  line-integral-exists:
  line-integral-exists F basis1 γ
  line-integral-exists F basis2 γ and
  basis-partition:

```

$basis1 \cup basis2 = basis$ $basis1 \cap basis2 = \{\}$
shows $line\text{-}integral\ F\ basis\ \gamma = (line\text{-}integral\ F\ basis1\ \gamma) + (line\text{-}integral\ F\ basis2\ \gamma)$
line-integral-exists F basis γ
apply (*simp add: line-integral-def*)
proof –
have 0: $integral\ \{0..1\}\ (\lambda x. (\sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + (\sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b))) = integral\ \{0..1\}\ (\lambda x. \sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + integral\ \{0..1\}\ (\lambda x. \sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b))$
using *Henstock-Kurzweil-Integration.integral-add line-integral-exists*
by (*auto simp add: line-integral-exists-def*)
have 1: $integral\ \{0..1\}\ (\lambda x. \sum b \in basis. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) = integral\ \{0..1\}\ (\lambda x. (\sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + (\sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)))$
by (*metis (mono-tags, lifting) basis-partition finite-Un finite-basis sum.union-disjoint*)
show $integral\ \{0..1\}\ (\lambda x. \sum b \in basis. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) = integral\ \{0..1\}\ (\lambda x. \sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + integral\ \{0..1\}\ (\lambda x. \sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b))$
using 0 1
by *auto*
have 2: $((\lambda x. (\sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + (\sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)))\ has\text{-}integral\ integral\ \{0..1\}\ (\lambda x. \sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + integral\ \{0..1\}\ (\lambda x. \sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)))\ \{0..1\}$
using *Henstock-Kurzweil-Integration.has-integral-add line-integral-exists has-integral-integral*
apply (*auto simp add: line-integral-exists-def*)
by *blast*
have 3: $(\lambda x. \sum b \in basis. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) = (\lambda x. (\sum b \in basis1. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)) + (\sum b \in basis2. F\ (\gamma\ x) \cdot b * (vector\text{-}derivative\ \gamma\ (at\ x\ within\ \{0..1\}) \cdot b)))$
by (*metis (mono-tags, lifting) basis-partition finite-Un finite-basis sum.union-disjoint*)

show *line-integral-exists* F *basis* γ
apply (*auto simp add: line-integral-exists-def has-integral-integral*)
using 2 3
using *has-integral-integrable-integral* **by** *fastforce*
qed

definition *common-boundary-sudivision-exists* **where**
common-boundary-sudivision-exists one-chain1 one-chain2 \equiv
 \exists *subdiv. chain-subdiv-chain one-chain1 subdiv* \wedge
 $\text{chain-subdiv-chain one-chain2 subdiv} \wedge$
 $(\forall (k, \gamma) \in \text{subdiv. valid-path } \gamma) \wedge$
 $\text{boundary-chain subdiv}$

lemma *common-boundary-sudivision-commutative*:
 $(\text{common-boundary-sudivision-exists one-chain1 one-chain2}) = (\text{common-boundary-sudivision-exists one-chain2 one-chain1})$
apply (*simp add: common-boundary-sudivision-exists-def*)
by *blast*

lemma *common-sudivision-imp-eq-line-integral*:
assumes (*common-boundary-sudivision-exists one-chain1 one-chain2*)
 $\text{boundary-chain one-chain1}$
 $\text{boundary-chain one-chain2}$
 $\forall (k, \gamma) \in \text{one-chain1. line-integral-exists } F \text{ basis } \gamma$
 finite one-chain1
 finite one-chain2
 finite basis
shows $\text{one-chain-line-integral } F \text{ basis one-chain1} = \text{one-chain-line-integral } F$
 basis one-chain2
 $\forall (k, \gamma) \in \text{one-chain2. line-integral-exists } F \text{ basis } \gamma$

proof –
obtain *subdiv* **where** *subdiv-props*:
 $\text{chain-subdiv-chain one-chain1 subdiv}$
 $\text{chain-subdiv-chain one-chain2 subdiv}$
 $(\forall (k, \gamma) \in \text{subdiv. valid-path } \gamma)$
 $(\text{boundary-chain subdiv})$
using *assms*
by (*auto simp add: common-boundary-sudivision-exists-def*)
have $i: \forall (k, \gamma) \in \text{subdiv. line-integral-exists } F \text{ basis } \gamma$
using $\text{one-chain-line-integral-eq-line-integral-on-sudivision}'(2)[OF \text{ subdiv-props}(1) \text{ assms}(2) \text{ subdiv-props}(4) \text{ assms}(4) \text{ subdiv-props}(3) \text{ assms}(5) \text{ assms}(7)]$
by *auto*
show $\text{one-chain-line-integral } F \text{ basis one-chain1} = \text{one-chain-line-integral } F \text{ basis one-chain2}$
using $\text{one-chain-line-integral-eq-line-integral-on-sudivision}'(1)[OF \text{ subdiv-props}(1) \text{ assms}(2) \text{ subdiv-props}(4) \text{ assms}(4) \text{ subdiv-props}(3) \text{ assms}(5) \text{ assms}(7)]$
 $\text{one-chain-line-integral-eq-line-integral-on-sudivision}(1)[OF \text{ subdiv-props}(2) \text{ assms}(3) \text{ subdiv-props}(4) i \text{ subdiv-props}(3) \text{ assms}(6) \text{ assms}(7)]$
by *auto*

show $\forall (k, \gamma) \in \text{one-chain2}. \text{line-integral-exists } F \text{ basis } \gamma$
using $\text{one-chain-line-integral-eq-line-integral-on-sudivision}(2)[OF \text{ subdiv-props}(2)$
 $\text{assms}(3) \text{ subdiv-props}(4) i \text{ subdiv-props}(3) \text{ assms}(6) \text{ assms}(7)]$
by *auto*
qed

definition *common-sudiv-exists where*

$\text{common-sudiv-exists one-chain1 one-chain2} \equiv$
 $\exists \text{ subdiv } ps1 \ ps2. \text{chain-sudiv-chain } (\text{one-chain1} - ps1) \text{ subdiv} \wedge$
 $\text{chain-sudiv-chain } (\text{one-chain2} - ps2) \text{ subdiv} \wedge$
 $(\forall (k, \gamma) \in \text{subdiv}. \text{valid-path } \gamma) \wedge$
 $(\text{boundary-chain subdiv}) \wedge$
 $(\forall (k, \gamma) \in ps1. \text{point-path } \gamma) \wedge$
 $(\forall (k, \gamma) \in ps2. \text{point-path } \gamma)$

lemma *common-sudiv-exists-comm:*

shows $\text{common-sudiv-exists } C1 \ C2 = \text{common-sudiv-exists } C2 \ C1$
by (*auto simp add: common-sudiv-exists-def*)

lemma *line-integral-degenerate-chain:*

assumes $(\forall (k, \gamma) \in \text{chain}. \text{point-path } \gamma)$
assumes *finite basis*
shows $\text{one-chain-line-integral } F \text{ basis chain} = 0$
proof (*simp add: one-chain-line-integral-def*)
have $\forall (k, g) \in \text{chain}. \text{line-integral } F \text{ basis } g = 0$
using *assms line-integral-point-path*
by *blast*
then have $\forall (k, g) \in \text{chain}. \text{real-of-int } k * \text{line-integral } F \text{ basis } g = 0$ **by** *auto*
then have $\bigwedge p. p \in \text{chain} \implies (\text{case } p \text{ of } (i, f) \Rightarrow \text{real-of-int } i * \text{line-integral } F$
 $\text{basis } f) = 0$
by *fastforce*
then show $(\sum x \in \text{chain}. \text{case } x \text{ of } (k, g) \Rightarrow \text{real-of-int } k * \text{line-integral } F \text{ basis } g) = 0$
by *simp*
qed

lemma *gen-common-sudiv-imp-common-sudiv:*

shows $(\text{common-sudiv-exists one-chain1 one-chain2}) = (\exists ps1 \ ps2. (\text{common-boundary-sudivision-exists}$
 $(\text{one-chain1} - ps1) (\text{one-chain2} - ps2)) \wedge (\forall (k, \gamma) \in ps1. \text{point-path } \gamma) \wedge (\forall (k,$
 $\gamma) \in ps2. \text{point-path } \gamma))$
by (*auto simp add: common-sudiv-exists-def common-boundary-sudivision-exists-def*)

lemma *common-sudiv-imp-gen-common-sudiv:*

assumes $(\text{common-boundary-sudivision-exists one-chain1 one-chain2})$
shows $(\text{common-sudiv-exists one-chain1 one-chain2})$
using *assms*
apply (*auto simp add: common-sudiv-exists-def common-boundary-sudivision-exists-def*)
by (*metis Diff-empty all-not-in-conv*)

lemma *one-chain-line-integral-point-paths*:
assumes *finite one-chain*
assumes *finite basis*
assumes $(\forall (k, \gamma) \in ps. \text{point-path } \gamma)$
shows *one-chain-line-integral F basis (one-chain - ps) = one-chain-line-integral F basis (one-chain)*
proof –
have $0: (\forall x \in ps. \text{case } x \text{ of } (k, g) \Rightarrow (\text{real-of-int } k * \text{line-integral F basis } g) = 0)$
using *line-integral-point-path assms*
by force
show *one-chain-line-integral F basis (one-chain - ps) = one-chain-line-integral F basis one-chain*
unfolding *one-chain-line-integral-def* **using** $0 \langle \text{finite one-chain} \rangle$
by (*force simp add: intro: comm-monoid-add-class.sum.mono-neutral-left*)
qed

lemma *boundary-chain-diff*:
assumes *boundary-chain one-chain*
shows *boundary-chain (one-chain - s)*
using *assms*
by (*auto simp add: boundary-chain-def*)

lemma *gen-common-subdivision-imp-eq-line-integral*:
assumes (*common-sudiv-exists one-chain1 one-chain2*)
boundary-chain one-chain1
boundary-chain one-chain2
 $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists F basis } \gamma$
finite one-chain1
finite one-chain2
finite basis
shows *one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis one-chain2*
 $\forall (k, \gamma) \in \text{one-chain2}. \text{line-integral-exists F basis } \gamma$
proof –
obtain *ps1 ps2* **where** *gen-sudiv: (common-boundary-subdivision-exists (one-chain1 - ps1) (one-chain2 - ps2))* $(\forall (k, \gamma) \in ps1. \text{point-path } \gamma)$ $(\forall (k, \gamma) \in ps2. \text{point-path } \gamma)$
using *assms(1) gen-common-sudiv-imp-common-sudiv*
by blast
show *one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis one-chain2*
using *one-chain-line-integral-point-paths gen-common-sudiv-imp-common-sudiv assms(2-7) gen-sudiv common-subdivision-imp-eq-line-integral(1)[OF gen-sudiv(1) boundary-chain-diff[OF assms(2)] boundary-chain-diff[OF assms(3)]]*
by auto
show $\forall (k, \gamma) \in \text{one-chain2}. \text{line-integral-exists F basis } \gamma$
proof –
obtain *sudiv* **where** *sudiv-props*:

```

    chain-subdiv-chain (one-chain1-ps1) subdiv
    chain-subdiv-chain (one-chain2-ps2) subdiv
    (∀ (k, γ) ∈ subdiv. valid-path γ)
    (boundary-chain subdiv)
    using gen-subdiv(1)
    by (auto simp add: common-boundary-sudivision-exists-def)
  have ∀ (k, γ) ∈ subdiv. line-integral-exists F basis γ
    using one-chain-line-integral-eq-line-integral-on-sudivision'(2)[OF subdiv-props(1)
    boundary-chain-diff[OF assms(2)] subdiv-props(4)] assms(4) subdiv-props(3) assms(5)
    assms(7)
    by blast
  then have i: ∀ (k, γ) ∈ one-chain2-ps2. line-integral-exists F basis γ
    using one-chain-line-integral-eq-line-integral-on-sudivision(2)[OF subdiv-props(2)
    boundary-chain-diff[OF assms(3)] subdiv-props(4)] subdiv-props(3) assms(6) assms(7)
    by blast
  then show ?thesis
    using gen-subdiv(3) line-integral-exists-point-path[OF assms(7)]
    by blast
qed
qed

```

```

lemma common-sudiv-exists-reft:
  assumes common-sudiv-exists C1 C2
  shows common-sudiv-exists C2 C1
  using assms
  apply (simp add: common-sudiv-exists-def)
  by auto

```

```

lemma chain-subdiv-path-singleton:
  shows chain-subdiv-path γ {(1, γ)}
proof -
  have rec-join [(1, γ)] = γ
    by (simp add: joinpaths-def)
  then have set [(1, γ)] = {(1, γ)} distinct [(1, γ)] rec-join [(1, γ)] = γ valid-chain-list
    [(1, γ)]
    by auto
  then show ?thesis
    by (metis (no-types) chain-subdiv-path.intros)
qed

```

```

lemma chain-subdiv-path-singleton-reverse:
  shows chain-subdiv-path (reversepath γ) {(-1, γ)}
proof -
  have rec-join [(-1, γ)] = reversepath γ
    by (simp add: joinpaths-def)
  then have set [(-1, γ)] = {(-1, γ)} distinct [(-1, γ)]
    rec-join [(-1, γ)] = reversepath γ valid-chain-list [(-1, γ)]
    by auto
  then have chain-subdiv-path (reversepath γ) (set [(-1, γ)])

```

using *chain-subdiv-path.intros* **by** *blast*
then show *?thesis*
by *simp*
qed

lemma *chain-subdiv-chain-refl*:
assumes *boundary-chain C*
shows *chain-subdiv-chain C C*
using *chain-subdiv-path-singleton chain-subdiv-path-singleton-reverse assms*
unfolding *chain-subdiv-chain-def boundary-chain-def pairwise-def* **using** *case-prodI2*
coeff-cube-to-path.simps
by (*rule-tac x= $\lambda x. \{x\}$ in exI*) *auto*

definition *reparam-weak* **where**

reparam-weak $\gamma 1 \gamma 2 \equiv \exists \varphi. (\forall x \in \{0..1\}. \gamma 1 x = (\gamma 2 \circ \varphi) x) \wedge \varphi$ *piecewise-C1-differentiable-on*
 $\{0..1\} \wedge \varphi(0) = 0 \wedge \varphi(1) = 1 \wedge \varphi^{-1} \{0..1\} = \{0..1\}$

definition *reparam* **where**

reparam $\gamma 1 \gamma 2 \equiv \exists \varphi. (\forall x \in \{0..1\}. \gamma 1 x = (\gamma 2 \circ \varphi) x) \wedge \varphi$ *piecewise-C1-differentiable-on*
 $\{0..1\} \wedge \varphi(0) = 0 \wedge \varphi(1) = 1 \wedge$ *bij-betw* $\varphi \{0..1\} \{0..1\} \wedge \varphi^{-1} \{0..1\} \subseteq \{0..1\}$
 $\wedge (\forall x \in \{0..1\}. \text{finite } (\varphi^{-1} \{x\}))$

lemma *reparam-weak-eq-refl*:

shows *reparam-weak* $\gamma 1 \gamma 1$
unfolding *reparam-weak-def*
apply (*rule-tac x=id in exI*)
by (*auto simp add: id-def piecewise-C1-differentiable-on-def C1-differentiable-on-def*
continuous-on-id)

lemma *line-integral-exists-smooth-one-base*:

assumes γ *C1-differentiable-on* $\{0..1\}$
continuous-on (*path-image* γ) ($\lambda x. F x \cdot b$)
shows *line-integral-exists* $F \{b\} \gamma$

proof –

have *gamma2-differentiable*: ($\forall x \in \{0 .. 1\}. \gamma$ *differentiable at* x)
using *assms(1)*
by (*auto simp add: valid-path-def C1-differentiable-on-eq*)
then have *gamma2-b-component-differentiable*: ($\forall x \in \{0 .. 1\}. (\lambda x. (\gamma x) \cdot b$)
differentiable at x)
by *auto*
then have ($\lambda x. (\gamma x) \cdot b$) *differentiable-on* $\{0..1\}$
using *differentiable-at-withinI*
by (*auto simp add: differentiable-on-def*)
then have *gama2-cont-comp*: *continuous-on* $\{0..1\}$ ($\lambda x. (\gamma x) \cdot b$)
using *differentiable-imp-continuous-on*
by *auto*
have *gamma2-cont:continuous-on* $\{0..1\} \gamma$

```

    using assms(1) C1-differentiable-imp-continuous-on
  by (auto simp add: valid-path-def)
  have iii: continuous-on {0..1} (λx. F (γ x) · b * (vector-derivative γ (at x within
{0..1}) · b))
  proof –
    have 0: continuous-on {0..1} (λx. F (γ x) · b)
      using assms(2) continuous-on-compose[OF gamma2-cont]
      by (auto simp add: path-image-def)
    obtain D where D: (∀ x∈{0..1}. (γ has-vector-derivative D x) (at x)) ∧
continuous-on {0..1} D
      using assms(1)
      by (auto simp add: C1-differentiable-on-def)
    then have *:∀ x∈{0..1}. vector-derivative γ (at x within{0..1}) = D x
      using vector-derivative-at vector-derivative-at-within-ivl
      by fastforce
    then have continuous-on {0..1} (λx. vector-derivative γ (at x within{0..1}))
      using continuous-on-eq D by force
    then have 1: continuous-on {0..1} (λx. (vector-derivative γ (at x within{0..1}))
· b)
      by(auto intro!: continuous-intros)
    show ?thesis
      using continuous-on-mult[OF 0 1] by auto
  qed
  then have (λx. F (γ x) · b * (vector-derivative γ (at x within {0..1}) · b))
integrable-on {0..1}
    using integrable-continuous-real
    by auto
  then show line-integral-exists F {b} γ
    by(auto simp add: line-integral-exists-def)
  qed

```

lemma *contour-integral-primitive-lemma:*

fixes $f :: \text{complex} \Rightarrow \text{complex}$ **and** $g :: \text{real} \Rightarrow \text{complex}$

assumes $a \leq b$

and $\bigwedge x. x \in s \implies (f \text{ has-field-derivative } f' x) (at x \text{ within } s)$

and $g \text{ piecewise-differentiable-on } \{a..b\} \bigwedge x. x \in \{a..b\} \implies g x \in s$

shows $((\lambda x. f'(g x) * \text{vector-derivative } g (at x \text{ within } \{a..b\}))$
 $\text{has-integral } (f(g b) - f(g a)) \{a..b\}$

proof –

obtain k **where** $k: \text{finite } k \forall x \in \{a..b\} - k. g \text{ differentiable } (at x \text{ within } \{a..b\})$

and $cg: \text{continuous-on } \{a..b\} g$

using assms **by** $(\text{auto simp: piecewise-differentiable-on-def})$

have $cfg: \text{continuous-on } \{a..b\} (\lambda x. f (g x))$

apply $(\text{rule continuous-on-compose } [OF cg, \text{unfolded o-def}])$

using assms

apply $(\text{metis field-differentiable-def field-differentiable-imp-continuous-at con-}$
 $\text{tinuous-on-eq-continuous-within continuous-on-subset image-subset-iff})$

done

{ **fix** $x::\text{real}$

```

assume  $a: a < x$  and  $b: x < b$  and  $xk: x \notin k$ 
then have  $g$  differentiable at  $x$  within  $\{a..b\}$ 
  using  $k$  by (simp add: differentiable-at-withinI)
then have ( $g$  has-vector-derivative vector-derivative  $g$  (at  $x$  within  $\{a..b\}$ )) (at
 $x$  within  $\{a..b\}$ )
  by (simp add: vector-derivative-works has-field-derivative-def scaleR-conv-of-real)
  then have  $gdiff: (g$  has-derivative  $(\lambda u. u * \text{vector-derivative } g \text{ (at } x \text{ within } \{a..b\}))$ 
 $\{a..b\}))$  (at  $x$  within  $\{a..b\}$ )
    by (simp add: has-vector-derivative-def scaleR-conv-of-real)
  have ( $f$  has-field-derivative  $(f' (g x))$ ) (at  $(g x)$  within  $g \text{ ' } \{a..b\}$ )
    using  $assms$  by (metis  $a$  atLeastAtMost-iff  $b$  DERIV-subset image-subset-iff
less-eq-real-def)
  then have  $fdiff: (f$  has-derivative  $(*) (f' (g x))$ ) (at  $(g x)$  within  $g \text{ ' } \{a..b\}$ )
    by (simp add: has-field-derivative-def)
  have  $((\lambda x. f (g x))$  has-vector-derivative  $f' (g x) * \text{vector-derivative } g \text{ (at } x \text{ within } \{a..b\})$ 
 $\{a..b\}))$  (at  $x$  within  $\{a..b\}$ )
    using diff-chain-within [OF  $gdiff$   $fdiff$ ]
    by (simp add: has-vector-derivative-def scaleR-conv-of-real o-def mult-ac)
} note  $*$  = this
show ?thesis
  apply (rule fundamental-theorem-of-calculus-interior-strong)
  using  $k$   $assms$   $cfg$   $*$ 
  apply (auto simp: at-within-Icc-at)
done

```

qed

lemma line-integral-primitive-lemma:

fixes $f :: 'a :: \{\text{euclidean-space, real-normed-field}\} \Rightarrow 'a :: \{\text{euclidean-space, real-normed-field}\}$
and

$g :: \text{real} \Rightarrow 'a$

assumes $\bigwedge(a :: 'a). a \in s \implies (f$ has-field-derivative $(f' a)$) (at a within s)

and g piecewise-differentiable-on $\{0..1\}$ $\bigwedge x. x \in \{0..1\} \implies g x \in s$

and $\text{base-vec} \in \text{Basis}$

shows $((\lambda x. ((f'(g x)) * (\text{vector-derivative } g \text{ (at } x \text{ within } \{0..1\}))) \cdot \text{base-vec})$
 $\text{has-integral } (((f(g 1)) \cdot \text{base-vec} - (f(g 0)) \cdot \text{base-vec})) \{0..1\}$

proof –

obtain k **where** $k: \text{finite } k \forall x \in \{0..1\} - k. g$ differentiable (at x within $\{0..1\}$)

and $cg: \text{continuous-on } \{0..1\}$ g

using $assms$ **by** (auto simp: piecewise-differentiable-on-def)

have $cfg: \text{continuous-on } \{0..1\}$ $(\lambda x. f (g x))$

apply (rule continuous-on-compose [OF cg , unfolded o-def])

using $assms$

apply (metis field-differentiable-def field-differentiable-imp-continuous-at con-
tinuous-on-eq-continuous-within continuous-on-subset image-subset-iff)

done

{ fix $x :: \text{real}$

assume $a: 0 < x$ **and** $b: x < 1$ **and** $xk: x \notin k$

then have g differentiable at x within $\{0..1\}$

using k **by** (simp add: differentiable-at-withinI)

then have (g has-vector-derivative vector-derivative g (at x within $\{0..1\}$)) (at x within $\{0..1\}$)
by (*simp add: vector-derivative-works has-field-derivative-def scaleR-conv-of-real*)
then have $gdiff$: (g has-derivative ($\lambda u.$ of-real $u * \text{vector-derivative } g$ (at x within $\{0..1\}$))) (at x within $\{0..1\}$)
by (*simp add: has-vector-derivative-def scaleR-conv-of-real*)
have (f has-field-derivative ($f' (g x)$)) (at $(g x)$ within $g^{-1} \{0..1\}$)
using *assms* **by** (*metis a atLeastAtMost-iff b DERIV-subset image-subset-iff less-eq-real-def*)
then have $fdiff$: (f has-derivative ($*$) ($f' (g x)$)) (at $(g x)$ within $g^{-1} \{0..1\}$)
by (*simp add: has-field-derivative-def*)
have ($(\lambda x. f (g x))$ has-vector-derivative $f' (g x) * \text{vector-derivative } g$ (at x within $\{0..1\}$)) (at x within $\{0..1\}$)
using *diff-chain-within [OF gdiff fdiff]*
by (*simp add: has-vector-derivative-def scaleR-conv-of-real o-def mult-ac*)
}
then have $*$: $\bigwedge x. x \in \{0 < .. < 1\} - k \implies ((\lambda x. f (g x)) \text{ has-vector-derivative } f' (g x) * \text{vector-derivative } g \text{ (at } x \text{ within } \{0..1\})) \text{ (at } x \text{ within } \{0..1\})$
by *auto*
have ($(\lambda x. ((f'(g x))) * ((\text{vector-derivative } g \text{ (at } x \text{ within } \{0..1\}))))$
 $\text{has-integral } (((f(g 1)) - (f(g 0)))) \{0..1\}$
using *fundamental-theorem-of-calculus-interior-strong[OF k(1) zero-le-one - cfg]*
using k *assms* cfg **by** (*auto simp: at-within-Icc-at*)
then have ($(\lambda x. (((f'(g x))) * ((\text{vector-derivative } g \text{ (at } x \text{ within } \{0..1\})))) \cdot$
 $\text{base-vec})$
 $\text{has-integral } (((f(g 1)) - (f(g 0)))) \cdot \text{base-vec} \{0..1\}$
using *has-integral-componentwise-iff assms(4)*
by *fastforce*
then show *?thesis* **using** *inner-mult-left'*
by (*simp add: inner-mult-left' inner-diff-left*)
qed

lemma *reparam-eq-line-integrals*:

assumes *reparam*: *reparam* $\gamma 1 \ \gamma 2$ **and**
pw-smooth: $\gamma 2$ *piecewise-C1-differentiable-on* $\{0..1\}$ **and**
cont: *continuous-on* (*path-image* $\gamma 2$) ($\lambda x. F x \cdot b$) **and**
line-integral-ex: *line-integral-exists* $F \{b\} \ \gamma 2$
shows *line-integral* $F \{b\} \ \gamma 1 = \text{line-integral } F \{b\} \ \gamma 2$
line-integral-exists $F \{b\} \ \gamma 1$

proof –

obtain φ **where** *phi*: ($\forall x \in \{0..1\}. \gamma 1 x = (\gamma 2 \circ \varphi) x$) φ *piecewise-C1-differentiable-on* $\{0..1\}$ $\varphi(0) = 0$ $\varphi(1) = 1$ *bij-betw* $\varphi \{0..1\} \{0..1\}$ $\varphi^{-1} \{0..1\} \subseteq \{0..1\} \forall x \in \{0..1\}$.
finite ($\varphi^{-1} \{x\}$)

using *reparam*

by (*auto simp add: reparam-def*)

obtain s **where** *s*: *finite* $s \ \varphi$ *C1-differentiable-on* $\{0..1\} - s$

using *phi*

by (*auto simp add: reparam-def piecewise-C1-differentiable-on-def*)

```

let ?s = s ∩ {0..1}
have s-inter: finite ?s φ C1-differentiable-on {0..1} - ?s
  using s
  apply blast
  by (metis Diff-Compl Diff-Diff-Int Diff-eq inf commute s(2))
have cont-phi: continuous-on {0..1} φ
  using phi
  by(auto simp add: reparam-def piecewise-C1-differentiable-on-imp-continuous-on)
obtain s' D where s'-D: finite s' (∀ x ∈ {0 .. 1} - s'. γ2 differentiable at x)
(∀ x ∈ {0..1} - s'. (γ2 has-vector-derivative D x) (at x)) ∧ continuous-on ({0..1}
- s') D
  using pw-smooth
  apply (auto simp add: valid-path-def piecewise-C1-differentiable-on-def C1-differentiable-on-eq)
  by (simp add: vector-derivative-works)
let ?s' = s' ∩ {0..1}
have gamma2-differentiable: finite ?s' (∀ x ∈ {0 .. 1} - ?s'. γ2 differentiable
at x) (∀ x ∈ {0..1} - ?s'. (γ2 has-vector-derivative D x) (at x)) ∧ continuous-on
({0..1} - ?s') D
  using s'-D
  apply blast
  using s'-D(2) apply auto[1]
  by (metis Diff-Int2 inf-top.left-neutral s'-D(3))
then have gamma2-b-component-differentiable: (∀ x ∈ {0 .. 1} - ?s'. (λx. (γ2
x) · b) differentiable at x)
  using differentiable-inner by force
then have (λx. (γ2 x) · b) differentiable-on {0..1} - ?s'
  using differentiable-at-withinI
  by (auto simp add: differentiable-on-def)
then have gama2-cont-comp: continuous-on ({0..1} - ?s') (λx. (γ2 x) · b)
  using differentiable-imp-continuous-on
  by auto

have s-in01: ?s ⊆ {0..1} by auto
have s'-in01: ?s' ⊆ {0..1} by auto
have phi-backimg-s': φ -' {0..1} ⊆ {0..1} using phi by auto
have inj-on φ {0..1} using phi(5) by (auto simp add: bij-betw-def)
have bij-phi: bij-betw φ {0..1} {0..1} using phi(5) by auto
have finite-bck-img-single: ∀ x ∈ {0..1}. finite (φ -' {x}) using phi by auto
then have finite-bck-img-single-s': ∀ x ∈ ?s'. finite (φ -' {x}) by auto
have gamma2-line-integrable: (λx. F (γ2 x) · b * (vector-derivative γ2 (at x
within {0..1}) · b)) integrable-on {0..1}
  using line-integral-ex
  by (simp add: line-integral-exists-def)

have finite-neg-img: finite (φ -' ?s')
  using finite-bck-img-single
  by (metis Int-iff finite-Int gamma2-differentiable(1) image-vimage-eq inf-img-fin-dom')
have gamma2-cont: continuous-on ({0..1} - ?s') γ2
  using gamma2-differentiable

```


by (*meson continuous-at-imp-continuous-on differentiable-imp-continuous-within*)
have *iii*: *continuous-on* ($\{0..1\} - ?s'$) ($\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2$
(at x within $\{0..1\}) \cdot b$)
proof –
have *0*: *continuous-on* ($\{0..1\} - ?s'$) ($\lambda x. F (\gamma 2 x) \cdot b$)
using *cont continuous-on-compose*[*OF gamma2-cont*] *continuous-on-compose2*
gamma2-cont
unfolding *path-image-def* **by** *fastforce*
have *D*: ($\forall x \in \{0..1\} - ?s'. (\gamma 2 \text{ has-vector-derivative } D x) (\text{at } x)) \wedge \text{continuous-on}$
 $(\{0..1\} - ?s') D$
using *gamma2-differentiable*
by *auto*
then have $*: \forall x \in \{0..1\} - ?s'. \text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\}) = D x$
using *vector-derivative-at vector-derivative-at-within-ivl*
by *fastforce*
then have *continuous-on* ($\{0..1\} - ?s'$) ($\lambda x. \text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\})$)
using *continuous-on-eq D*
by *metis*
then have *1*: *continuous-on* ($\{0..1\} - ?s'$) ($\lambda x. (\text{vector-derivative } \gamma 2 (\text{at } x$
within $\{0..1\})) \cdot b$)
by (*auto intro!*: *continuous-intros*)
show *?thesis*
using *continuous-on-mult*[*OF 0 1*]
by *auto*
qed
have *iv*: $\varphi(0) \leq \varphi(1)$
using *phi(3) phi(4)*
by (*simp add: reparam-def*)
have *v*: $\varphi \{0..1\} \subseteq \{0..1\}$
using *phi*
by (*auto simp add: reparam-def bij-betw-def*)
obtain *D* **where** *D-props*: ($\forall x \in \{0..1\} - ?s. (\varphi \text{ has-vector-derivative } D x) (\text{at}$
x)
using *s*
by (*auto simp add: C1-differentiable-on-def*)
then have ($\bigwedge x. x \in (\{0..1\} - ?s) \implies (\varphi \text{ has-vector-derivative } D x) (\text{at } x \text{ within}$
 $\{0..1\})$)
using *has-vector-derivative-at-within*
by *blast*
then have *vi*: ($\bigwedge x. x \in (\{0..1\} - ?s) \implies (\varphi \text{ has-real-derivative } D x) (\text{at } x$
within $\{0..1\})$)
using *has-real-derivative-iff-has-vector-derivative*
by *blast*
have *a*: ($(\lambda x. D x * (F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } (\varphi x) \text{ within}$
 $\{0..1\}) \cdot b))$) *has-integral*
 $\text{integral } \{\varphi 0.. \varphi 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x$
within $\{0..1\}) \cdot b)$)
 $(\{0..1\})$
proof –

have a : $\text{integral } \{\varphi \ 1.. \varphi \ 0\} (\lambda x. F (\gamma 2 \ x) \cdot b * (\text{vector-derivative } \gamma 2 \ (\text{at } x \ \text{within } \{0..1\}) \cdot b)) = 0$ **using** $\text{integral-singleton integral-empty iv}$
by $(\text{simp add: phi(3) phi(4)})$
have b : $((\lambda x. D \ x *_R (F (\gamma 2 (\varphi \ x)) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } (\varphi \ x) \ \text{within } \{0..1\}) \cdot b))) \ \text{has-integral}$
 $\text{integral } \{\varphi \ 0.. \varphi \ 1\} (\lambda x. F (\gamma 2 \ x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \ \text{within } \{0..1\}) \cdot b)) - \text{integral } \{\varphi \ 1.. \varphi \ 0\} (\lambda x. F (\gamma 2 \ x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \ \text{within } \{0..1\}) \cdot b)))$
 $\{0..1\}$
apply $(\text{rule has-integral-substitution-general-}[OF \ s\text{-inter}(1) \ \text{zero-le-one gamma2-differentiable}(1) \ v \ \text{gamma2-line-integrable iii cont-phi finite-bck-img-single-s'}])$
proof–
have $\text{surj-on } \{0..1\} \ \varphi$
using bij-phi
by $(\text{metis (full-types) bij-betw-def image-subsetI rangeI})$
then show $\text{surj-on } ?s' \ \varphi$ **using** bij-phi s'-in01
by blast
show $\text{inj-on } \varphi (\{0..1\} \cup (?s \cup \varphi -' ?s'))$
proof–
have i : $\text{inj-on } \varphi \ \{0..1\}$ **using** bij-phi
using bij-betw-def **by** blast
have ii : $(\{0..1\} \cup (?s \cup \varphi -' ?s')) = \{0..1\}$ **using** $\text{phi-backimg-s' s-in01}$
 $s'\text{-in01}$
by blast
show $?thesis$ **using** $i \ ii$ **by** auto
qed
show $\bigwedge x. x \in \{0..1\} - ?s \implies (\varphi \ \text{has-real-derivative } D \ x) (\text{at } x \ \text{within } \{0..1\})$
using vi **by** blast
qed
show $?thesis$ **using** $a \ b$ **by** auto
qed
then have b : $\text{integral } \{0..1\} (\lambda x. D \ x * (F (\gamma 2 (\varphi \ x)) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } (\varphi \ x) \ \text{within } \{0..1\}) \cdot b))) =$
 $\text{integral } \{\varphi \ 0.. \varphi \ 1\} (\lambda x. F (\gamma 2 \ x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \ \text{within } \{0..1\}) \cdot b))$
by auto
have $\text{gamma2-vec-diffable}$: $\bigwedge x::\text{real}. x \in \{0..1\} - ((\varphi -' ?s') \cup ?s) \implies ((\gamma 2 \ o \ \varphi) \ \text{has-vector-derivative vector-derivative } (\gamma 2 \ o \ \varphi) (\text{at } x)) (\text{at } x)$
proof–
fix $x::\text{real}$
assume ass : $x \in \{0..1\} - ((\varphi -' ?s') \cup ?s)$
have zer-le-x-le-1 : $0 \leq x \wedge x \leq 1$ **using** ass
by simp
show $((\gamma 2 \ o \ \varphi) \ \text{has-vector-derivative vector-derivative } (\gamma 2 \ o \ \varphi) (\text{at } x)) (\text{at } x)$
proof–
have $**$: $\gamma 2$ *differentiable* *at* $(\varphi \ x)$
using $\text{gamma2-differentiable}(2) \ ass \ v$
by blast
have $***$: φ *differentiable* *at* x

```

    using s ass
    by (auto simp add: C1-differentiable-on-eq)
  then show (( $\gamma 2 \circ \varphi$ ) has-vector-derivative vector-derivative ( $\gamma 2 \circ \varphi$ ) (at x))
(at x)
    using differentiable-chain-at[OF *** **]
    by (auto simp add: vector-derivative-works)
  qed
  qed
  then have gamma2-vec-deriv-within:  $\bigwedge x::real. x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s) \implies$ 
vector-derivative ( $\gamma 2 \circ \varphi$ ) (at x) = vector-derivative ( $\gamma 2 \circ \varphi$ ) (at x within
 $\{0..1\}$ )
    using vector-derivative-at-within-ivl[OF gamma2-vec-diffable]
    by auto
  have  $\forall x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s). D x * (vector-derivative \gamma 2 (at (\varphi x)$ 
within  $\{0..1\}) \cdot b) = (vector-derivative (\gamma 2 \circ \varphi) (at x within \{0..1\}) \cdot b)$ 
  proof
    fix x::real
    assume ass:  $x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s)$ 
    then have 0:  $\varphi$  differentiable (at x)
    using s by (auto simp add: C1-differentiable-on-def differentiable-def has-vector-derivative-def)
    obtain D2 where ( $\varphi$  has-vector-derivative D2) (at x)
    using D-props ass by blast
    have  $\varphi x \in \{0..1\} - ?s'$ 
    using phi(5) ass by (metis Diff-Un Diff-iff Int-iff bij-betw-def image-eqI
vimageI)
    then have 1:  $\gamma 2$  differentiable (at ( $\varphi x$ ))
    using gamma2-differentiable
    by auto
    have 3: vector-derivative  $\gamma 2$  (at ( $\varphi x$ )) = vector-derivative  $\gamma 2$  (at ( $\varphi x$ ) within
 $\{0..1\}$ )
    proof-
      have *:  $0 \leq \varphi x \wedge \varphi x \leq 1$  using phi(5) ass
      using  $\langle \varphi x \in \{0..1\} - s' \cap \{0..1\} \rangle$  by auto
      then have **: ( $\gamma 2$  has-vector-derivative (vector-derivative  $\gamma 2$  (at ( $\varphi x$ )))) (at
( $\varphi x$ ))
      using 1 vector-derivative-works by auto
      show ?thesis
      using * vector-derivative-at-within-ivl[OF **] by auto
    qed
    show  $D x * (vector-derivative \gamma 2 (at (\varphi x) within \{0..1\}) \cdot b) = vector-derivative$ 
( $\gamma 2 \circ \varphi$ ) (at x within  $\{0..1\}) \cdot b$ 
    using vector-derivative-chain-at[OF 0 1]
    apply (auto simp add: gamma2-vec-deriv-within[OF ass, symmetric] 3[symmetric])
    using D-props ass vector-derivative-at
    by fastforce
  qed
  then have c:  $\bigwedge x. x \in (\{0..1\} - ((\varphi - ' ?s') \cup ?s)) \implies D x * (F (\gamma 2 (\varphi x)) \cdot b$ 
* (vector-derivative  $\gamma 2$  (at ( $\varphi x$ ) within  $\{0..1\}) \cdot b) =$ 
 $F (\gamma 2 (\varphi x)) \cdot b * (vector-derivative (\gamma 2 \circ \varphi) (at x$ 

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within $\{0..1\} \cdot b$
by auto
then have d : *integral* $\{0..1\}$ $(\lambda x. D x * (F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } (\varphi x) \text{ within } \{0..1\}) \cdot b))) =$
 $\text{integral } \{0..1\} (\lambda x. F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) (\text{at } x \text{ within } \{0..1\}) \cdot b))$
proof–
have *negligible* $((\varphi - ' ?s') \cup ?s)$ **using** *finite-neg-img s(1)* **by auto**
then show *?thesis*
using *c integral-spike* **by** $(\text{metis}(\text{no-types}, \text{lifting}))$
qed
have *phi-in-int*: $(\bigwedge x. x \in \{0..1\} \implies \varphi x \in \{0..1\})$ **using** *phi*
using *v* **by** *blast*
then have e : $((\lambda x. F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) (\text{at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } \text{integral } \{\varphi 0.. \varphi 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\}) \cdot b))) \{0..1\}$
proof–
have *negligible* $?s$ **using** *s-inter(1)* **by auto**
have 0 : *negligible* $((\varphi - ' ?s') \cup ?s)$ **using** *finite-neg-img s(1)* **by auto**
have c' : $\forall x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s). D x * (F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } (\varphi x) \text{ within } \{0..1\}) \cdot b)) =$
 $F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) (\text{at } x \text{ within } \{0..1\}) \cdot b)$
using *c* **by** *blast*
have *has-integral-spike-eq'*: $\bigwedge s t f g y. \text{negligible } s \implies \forall x \in t - s. g x = f x \implies (f \text{ has-integral } y) t = (g \text{ has-integral } y) t$
using *has-integral-spike-eq* **by** *metis*
show *?thesis*
using *a has-integral-spike-eq'[OF 0 c']* **by** *blast*
qed
then have f : $((\lambda x. F (\gamma 1 x) \cdot b * (\text{vector-derivative } \gamma 1 (\text{at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } \text{integral } \{\varphi 0.. \varphi 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\}) \cdot b))) \{0..1\}$
proof–
assume *ass*: $((\lambda x. F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) (\text{at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } \text{integral } \{\varphi 0.. \varphi 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\}) \cdot b))) \{0..1\}$
have $*$: $\forall x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s) \cup \{0,1\}. (\lambda x. F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) (\text{at } x \text{ within } \{0..1\}) \cdot b)) x =$
 $(\lambda x. F (\gamma 1 x) \cdot b * (\text{vector-derivative } \gamma 1 (\text{at } x \text{ within } \{0..1\}) \cdot b)) x$
proof–
have $\forall x \in \{0 <.. < 1\} - (\varphi - ' ?s' \cup ?s). (\text{vector-derivative } (\gamma 2 \circ \varphi) (\text{at } x$

$\text{within } \{0..1\} \cdot b) = (\text{vector-derivative } (\gamma 1) \text{ (at } x \text{ within } \{0..1\}) \cdot b)$
proof
have i : $\text{open } (\{0 < .. < 1\} - ((\varphi - ' ?s') \cup ?s))$ **using** $\text{open-diff } s(1)$
 $\text{open-greaterThanLessThan finite-neg-img}$
by $(\text{simp add: open-diff})$
have ii : $\forall x \in \{0 < .. < 1 :: \text{real}\} - (\varphi - ' ?s' \cup ?s). (\gamma 2 \circ \varphi) x = \gamma 1 x$ **using**
 $\text{phi}(1)$
by auto
fix $x :: \text{real}$
assume ass : $x \in \{0 < .. < 1 :: \text{real}\} - ((\varphi - ' ?s') \cup ?s)$
then have iii : $(\gamma 2 \circ \varphi \text{ has-vector-derivative vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\})) \text{ (at } x)$
by $(\text{metis (no-types) Diff-iff add.commute add-strict-mono ass atLeast-AtMost-iff gamma2-vec-deriv-within gamma2-vec-diffable greaterThanLessThan-iff less-irrefl not-le})$

then have iv : $(\gamma 1 \text{ has-vector-derivative vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\})) \text{ (at } x)$
using $\text{has-derivative-transform-within-open } i \text{ ii ass}$
apply $(\text{auto simp add: has-vector-derivative-def})$
apply $(\text{meson ass has-derivative-transform-within-open } i \text{ ii})$
apply $(\text{meson ass has-derivative-transform-within-open } i \text{ ii})$
by $(\text{meson ass has-derivative-transform-within-open } i \text{ ii})$
have v : $0 \leq x \leq 1$ **using** ass **by** auto
have 0 : $\text{vector-derivative } \gamma 1 \text{ (at } x \text{ within } \{0..1\}) = \text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\})$
using $\text{vector-derivative-at-within-ivl}[OF \text{ iv } v(1) v(2) \text{ zero-less-one}]$
by force
have 1 : $\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) = \text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\})$
using $\text{vector-derivative-at-within-ivl}[OF \text{ iii } v(1) v(2) \text{ zero-less-one}]$
by force
then have $\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) = \text{vector-derivative } \gamma 1 \text{ (at } x \text{ within } \{0..1\})$
using $0 \ 1$ **by** auto
then show $\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b = \text{vector-derivative } \gamma 1 \text{ (at } x \text{ within } \{0..1\}) \cdot b$ **by** auto
qed
then have i : $\forall x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s) \cup \{0, 1\}. (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b) = (\text{vector-derivative } (\gamma 1) \text{ (at } x \text{ within } \{0..1\}) \cdot b)$
by auto
have ii : $\forall x \in \{0..1\} - ((\varphi - ' ?s') \cup ?s) \cup \{0, 1\}. F (\gamma 1 x) \cdot b = F (\gamma 2 (\varphi x)) \cdot b$
using $\text{phi}(1)$
by auto
show $?thesis$
using $i \text{ ii}$ **by** metis
qed

have **: *negligible* $((\varphi - ' ?s') \cup ?s \cup \{0, 1\})$ **using** $s(1)$ *finite-neg-img* **by**
auto
have *has-integral-spike-eq'*: $\bigwedge s t g f y. \text{negligible } s \implies$
 $\forall x \in t - s. g \ x = f \ x \implies (f \text{ has-integral } y) \ t = (g$
has-integral $y) \ t$
using *has-integral-spike-eq* **by** *metis*
show *?thesis*
using *has-integral-spike-eq'* $[OF \ ** \ *]$ *ass*
by *blast*
qed
then show *line-integral-exists* $F \ \{b\} \ \gamma 1$
using *phi* **by** *(auto simp add: line-integral-exists-def)*
have *integral* $(\{0..1\}) \ (\lambda x. F \ (\gamma 2 \ (\varphi \ x)) \cdot b \ * \ (\text{vector-derivative} \ (\gamma 2 \circ \varphi) \ (\text{at } x$
within $\{0..1\}) \cdot b)) =$
 $\text{integral} \ (\{0..1\}) \ (\lambda x. F \ (\gamma 1 \ x) \cdot b \ * \ (\text{vector-derivative} \ \gamma 1 \ (\text{at } x$
within $\{0..1\}) \cdot b))$
using *integral-unique* $[OF \ e]$ *integral-unique* $[OF \ f]$
by *metis*
moreover have *integral* $(\{0..1\}) \ (\lambda x. F \ (\gamma 2 \ (\varphi \ x)) \cdot b \ * \ (\text{vector-derivative} \ (\gamma 2$
 $\circ \varphi) \ (\text{at } x \text{ within } \{0..1\}) \cdot b)) =$
 $\text{integral} \ (\{0..1\}) \ (\lambda x. F \ (\gamma 2 \ x) \cdot b \ * \ (\text{vector-derivative} \ \gamma 2 \ (\text{at}$
x within $\{0..1\}) \cdot b))$
using *b d phi* **by** *(auto simp add:)*
ultimately show *line-integral* $F \ \{b\} \ \gamma 1 = \text{line-integral } F \ \{b\} \ \gamma 2$
using *phi* **by** *(auto simp add: line-integral-def)*
qed

lemma *reparam-weak-eq-line-integrals:*

assumes *reparam-weak* $\gamma 1 \ \gamma 2$
 $\gamma 2$ *C1-differentiable-on* $\{0..1\}$
continuous-on $(\text{path-image } \gamma 2) \ (\lambda x. F \ x \cdot b)$
shows *line-integral* $F \ \{b\} \ \gamma 1 = \text{line-integral } F \ \{b\} \ \gamma 2$
line-integral-exists $F \ \{b\} \ \gamma 1$
proof –
obtain φ **where** *phi*: $(\forall x \in \{0..1\}. \gamma 1 \ x = (\gamma 2 \circ \varphi) \ x)$ φ *piecewise-C1-differentiable-on*
 $\{0..1\}$ $\varphi(0) = 0$ $\varphi(1) = 1$ $\varphi' \ \{0..1\} = \{0..1\}$
using *assms(1)*
by *(auto simp add: reparam-weak-def)*
obtain s **where** s : *finite* $s \ \varphi$ *C1-differentiable-on* $\{0..1\} - s$
using *phi*
by *(auto simp add: reparam-weak-def piecewise-C1-differentiable-on-def)*
have *cont-phi*: *continuous-on* $\{0..1\} \ \varphi$
using *phi*
by *(auto simp add: reparam-weak-def piecewise-C1-differentiable-on-imp-continuous-on)*
have *gamma2-differentiable*: $(\forall x \in \{0 .. 1\}. \gamma 2 \ \text{differentiable at } x)$
using *assms(2)*
by *(auto simp add: valid-path-def C1-differentiable-on-eq)*
then have *gamma2-b-component-differentiable*: $(\forall x \in \{0 .. 1\}. (\lambda x. (\gamma 2 \ x) \cdot b)$
differentiable at $x)$

```

  by auto
then have (λx. (γ2 x) · b) differentiable-on {0..1}
  using differentiable-at-withinI
  by (auto simp add: differentiable-on-def)
then have gama2-cont-comp: continuous-on {0..1} (λx. (γ2 x) · b)
  using differentiable-imp-continuous-on
  by auto
have gamma2-cont:continuous-on {0..1} γ2
  using assms(2) C1-differentiable-imp-continuous-on
  by (auto simp add: valid-path-def)
have iii: continuous-on {0..1} (λx. F (γ2 x) · b * (vector-derivative γ2 (at x
within {0..1}) · b))
proof-
  have 0: continuous-on {0..1} (λx. F (γ2 x) · b)
    using assms(3) continuous-on-compose[OF gamma2-cont]
    by (auto simp add: path-image-def)
  obtain D where D: (∀ x∈{0..1}. (γ2 has-vector-derivative D x) (at x)) ∧
continuous-on {0..1} D
    using assms(2) by (auto simp add: C1-differentiable-on-def)
  then have *:∀ x∈{0..1}. vector-derivative γ2 (at x within{0..1}) = D x
    using vector-derivative-at vector-derivative-at-within-ivl
    by fastforce
  then have continuous-on {0..1} (λx. vector-derivative γ2 (at x within{0..1}))
    using continuous-on-eq D by force
  then have 1: continuous-on {0..1} (λx. (vector-derivative γ2 (at x within{0..1}))
· b)
    by (auto intro!: continuous-intros)
  show ?thesis
    using continuous-on-mult[OF 0 1] by auto
qed
have iv: φ(0) ≤ φ(1)
  using phi(3) phi(4) by (simp add: reparam-weak-def)
have v: φ{0..1} ⊆ {0..1}
  using phi(5) by (simp add: reparam-weak-def)
obtain D where D-props: (∀ x∈{0..1} - s. (φ has-vector-derivative D x) (at x))
  using s
  by (auto simp add: C1-differentiable-on-def)
then have (∧x. x ∈ ({0..1} - s) ⇒ (φ has-vector-derivative D x) (at x within
{0..1}))
  using has-vector-derivative-at-within by blast
then have vi: (∧x. x ∈ ({0..1} - s) ⇒ (φ has-real-derivative D x) (at x within
{0..1}))
  using has-real-derivative-iff-has-vector-derivative
  by blast
have a:((λx. D x * (F (γ2 (φ x)) · b * (vector-derivative γ2 (at (φ x) within
{0..1}) · b))) has-integral
  integral {φ 0..φ 1} (λx. F (γ2 x) · b * (vector-derivative γ2 (at x
within {0..1}) · b)))
  {0..1}

```

```

    using has-integral-substitution-strong[OF s(1) zero-le-one iv v iii cont-phi vi]
    by simp
  then have b: integral {0..1} (λx. D x * (F (γ2 (φ x)) · b * (vector-derivative
γ2 (at (φ x) within {0..1}) · b))) =
      integral {φ 0..φ 1} (λx. F (γ2 x) · b * (vector-derivative γ2 (at
x within {0..1}) · b))
    by auto
  have gamma2-vec-diffable: ∧x::real. x ∈ {0..1} - s ⇒ ((γ2 o φ) has-vector-derivative
vector-derivative (γ2 o φ) (at x)) (at x)
  proof -
    fix x::real
    assume ass: x ∈ {0..1} - s
    have zer-le-x-le-1: 0 ≤ x ∧ x ≤ 1 using ass by auto
    show ((γ2 o φ) has-vector-derivative vector-derivative (γ2 o φ) (at x)) (at x)
    proof -
      have **: γ2 differentiable at (φ x)
        using phi gamma2-differentiable
        by (auto simp add: zer-le-x-le-1)
      have ***: φ differentiable at x
        using s ass
        by (auto simp add: C1-differentiable-on-eq)
      then show ((γ2 o φ) has-vector-derivative vector-derivative (γ2 o φ) (at x))
(at x)
        using differentiable-chain-at[OF *** **]
        by (auto simp add: vector-derivative-works)
    qed
  qed
  then have gamma2-vec-deriv-within: ∧x::real. x ∈ {0..1} - s ⇒ vector-derivative
(γ2 o φ) (at x) = vector-derivative (γ2 o φ) (at x within {0..1})
    using vector-derivative-at-within-ivl[OF gamma2-vec-diffable]
    by auto
  have ∀x∈{0..1} - s. D x * (vector-derivative γ2 (at (φ x) within {0..1}) · b)
= (vector-derivative (γ2 o φ) (at x within {0..1}) · b)
  proof
    fix x::real
    assume ass: x ∈ {0..1} - s
    then have 0: φ differentiable (at x)
      using s
    by (auto simp add: C1-differentiable-on-def differentiable-def has-vector-derivative-def)
    obtain D2 where (φ has-vector-derivative D2) (at x)
      using D-props ass
      by blast
    have φ x ∈ {0..1}
      using phi(5) ass
    by (auto simp add: reparam-weak-def)
    then have 1: γ2 differentiable (at (φ x))
      using gamma2-differentiable
    by auto
    have 3: vector-derivative γ2 (at (φ x)) = vector-derivative γ2 (at (φ x) within

```



```

{0..1})
  proof-
    have *:  $0 \leq \varphi x \wedge \varphi x \leq 1$  using phi(5) ass by auto
    then have **:  $(\gamma 2 \text{ has-vector-derivative } (\text{vector-derivative } \gamma 2 \text{ (at } (\varphi x)))) \text{ (at } (\varphi x))$ 
      using 1 vector-derivative-works
      by auto
    show ?thesis
      using * vector-derivative-at-within-ivl[OF **]
      by auto
    qed
    show  $D x * (\text{vector-derivative } \gamma 2 \text{ (at } (\varphi x) \text{ within } \{0..1\}) \cdot b) = \text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b$ 
      using vector-derivative-chain-at[OF 0 1]
      apply (auto simp add: gamma2-vec-deriv-within[OF ass, symmetric] 3[symmetric])
      using D-props ass vector-derivative-at
      by fastforce
    qed
    then have c:  $\bigwedge x. x \in (\{0..1\} - s) \implies D x * (F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } \gamma 2 \text{ (at } (\varphi x) \text{ within } \{0..1\}) \cdot b)) = F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b)$ 
      by auto
    then have d:  $\text{integral } (\{0..1\}) (\lambda x. D x * (F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } \gamma 2 \text{ (at } (\varphi x) \text{ within } \{0..1\}) \cdot b))) = \text{integral } (\{0..1\}) (\lambda x. F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b))$ 
      proof-
        have negligible s using s(1) by auto
        then show ?thesis
          using c integral-spike by (metis(no-types, lifting))
        qed
        have phi-in-int:  $(\bigwedge x. x \in \{0..1\} \implies \varphi x \in \{0..1\})$  using phi
          by(auto simp add:)
        then have e:  $((\lambda x. F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b)) \text{ has-integral } \text{integral } \{\varphi 0.. \varphi 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 \text{ (at } x \text{ within } \{0..1\}) \cdot b))) \{0..1\}$ 
          proof-
            have 0:negligible s using s(1) by auto
            have c':  $\forall x \in \{0..1\} - s. D x * (F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } \gamma 2 \text{ (at } (\varphi x) \text{ within } \{0..1\}) \cdot b)) = F (\gamma 2 (\varphi x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b)$ 
              using c by auto
            have has-integral-spike-eq':  $\bigwedge s t f g y. \text{negligible } s \implies \forall x \in t - s. g x = f x \implies (f \text{ has-integral } y) t = (g \text{ has-integral } y) t$ 
              using has-integral-spike-eq by metis
          qed
        then

```

show *?thesis*
using *a has-integral-spike-eq'[OF 0 c]* **by** *blast*
qed
then have *f: ((λx. F (γ1 x) · b * (vector-derivative γ1 (at x within {0..1}) · b)) has-integral*

$$\int_{\{0..1\}} \{\varphi \ 0.. \varphi \ 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\}) \cdot b))$$

proof–
assume *ass: ((λx. F (γ2 (φ x)) · b * (vector-derivative (γ2 ∘ φ) (at x within {0..1}) · b)) has-integral*

$$\int_{\{0..1\}} \{\varphi \ 0.. \varphi \ 1\} (\lambda x. F (\gamma 2 x) \cdot b * (\text{vector-derivative } \gamma 2 (\text{at } x \text{ within } \{0..1\}) \cdot b))$$

have **: ∀ x ∈ {0..1} – (s ∪ {0,1}). (λx. F (γ2 (φ x)) · b * (vector-derivative (γ2 ∘ φ) (at x within {0..1}) · b)) x =*

$$(\lambda x. F (\gamma 1 x) \cdot b * (\text{vector-derivative } \gamma 1 (\text{at } x \text{ within } \{0..1\}) \cdot b)) x$$

proof–
have *∀ x ∈ {0 <.. < 1} – s. (vector-derivative (γ2 ∘ φ) (at x within {0..1}) · b) = (vector-derivative (γ1) (at x within {0..1}) · b)*
proof
have *i: open ({0 <.. < 1} – s) using open-diff s open-greaterThanLessThan*
by *blast*
have *ii: ∀ x ∈ {0 <.. < 1 :: real} – s. (γ2 ∘ φ) x = γ1 x using phi(1)*
by *auto*
fix *x :: real*
assume *ass: x ∈ {0 <.. < 1 :: real} – s*
then have *iii: (γ2 ∘ φ has-vector-derivative vector-derivative (γ2 ∘ φ) (at x within {0..1})) (at x)*
using *has-vector-derivative-at-within gamma2-vec-deriv-within gamma2-vec-diffable*
by *auto*

then have *iv: (γ1 has-vector-derivative vector-derivative (γ2 ∘ φ) (at x within {0..1})) (at x)*
using *has-derivative-transform-within-open i ii ass*
apply *(auto simp add: has-vector-derivative-def)*
by *force*
have *v: 0 ≤ x x ≤ 1 using ass by auto*
have *0: vector-derivative γ1 (at x within {0..1}) = vector-derivative (γ2 ∘ φ) (at x within {0..1})*
using *vector-derivative-at-within-ivl[OF iv v(1) v(2) zero-less-one]*
by *force*
have *1: vector-derivative (γ2 ∘ φ) (at x within {0..1}) = vector-derivative (γ2 ∘ φ) (at x within {0..1})*
using *vector-derivative-at-within-ivl[OF iii v(1) v(2) zero-less-one]*
by *force*
then have *vector-derivative (γ2 ∘ φ) (at x within {0..1}) = vector-derivative γ1 (at x within {0..1})*

using $0\ 1$ **by** *auto*
then show $\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b = \text{vector-derivative } \gamma 1 \text{ (at } x \text{ within } \{0..1\}) \cdot b$ **by** *auto*
qed
then have $i: \forall x \in \{0..1\} - (s \cup \{0,1\}). (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b) = (\text{vector-derivative } \gamma 1 \text{ (at } x \text{ within } \{0..1\}) \cdot b)$
by *auto*
have $ii: \forall x \in \{0..1\} - (s \cup \{0,1\}). F (\gamma 1\ x) \cdot b = F (\gamma 2 (\varphi\ x)) \cdot b$
using $\text{phi}(1)$ **by** *auto*
show *?thesis*
using $i\ ii$ **by** *auto*
qed
have $**:$ *negligible* $(s \cup \{0,1\})$ **using** $s(1)$ **by** *auto*
have $\text{has-integral-spike-eq}': \bigwedge s\ t\ g\ f\ y. \text{negligible } s \implies \forall x \in t - s. g\ x = f\ x \implies (f \text{ has-integral } y)\ t = (g \text{ has-integral } y)\ t$
using $\text{has-integral-spike-eq}$ **by** *metis*
show *?thesis*
using $\text{has-integral-spike-eq}'[OF\ **\ *]$ *ass*
by *blast*
qed
then show $\text{line-integral-exists } F\ \{b\}\ \gamma 1$
using phi **by** $(\text{auto simp add: line-integral-exists-def})$
have $\text{integral } (\{0..1\}) (\lambda x. F (\gamma 2 (\varphi\ x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b)) = \text{integral } (\{0..1\}) (\lambda x. F (\gamma 1\ x) \cdot b * (\text{vector-derivative } \gamma 1 \text{ (at } x \text{ within } \{0..1\}) \cdot b))$
using $\text{integral-unique}[OF\ e]\ \text{integral-unique}[OF\ f]$
by *metis*
moreover have $\text{integral } (\{0..1\}) (\lambda x. F (\gamma 2 (\varphi\ x)) \cdot b * (\text{vector-derivative } (\gamma 2 \circ \varphi) \text{ (at } x \text{ within } \{0..1\}) \cdot b)) = \text{integral } (\{0..1\}) (\lambda x. F (\gamma 2\ x) \cdot b * (\text{vector-derivative } \gamma 2 \text{ (at } x \text{ within } \{0..1\}) \cdot b))$
using $b\ d\ \text{phi}$ **by** (auto simp add:)
ultimately show $\text{line-integral } F\ \{b\}\ \gamma 1 = \text{line-integral } F\ \{b\}\ \gamma 2$
using phi **by** $(\text{auto simp add: line-integral-def})$
qed

lemma $\text{line-integral-sum-basis}$:
assumes $\text{finite } (basis::('a::\text{euclidean-space})\ \text{set})\ \forall b \in \text{basis}. \text{line-integral-exists } F\ \{b\}\ \gamma$
shows $\text{line-integral } F\ \text{basis}\ \gamma = (\sum b \in \text{basis}. \text{line-integral } F\ \{b\}\ \gamma)$
 $\text{line-integral-exists } F\ \text{basis}\ \gamma$
using *assms*
proof (induction basis)
show $\text{line-integral } F\ \{\}\ \gamma = (\sum b \in \{\}. \text{line-integral } F\ \{b\}\ \gamma)$
by $(\text{auto simp add: line-integral-def})$
show $\forall b \in \{\}. \text{line-integral-exists } F\ \{b\}\ \gamma \implies \text{line-integral-exists } F\ \{\}\ \gamma$
by $(\text{simp add: line-integral-exists-def integrable-0})$

next
fix *basis*::('a::euclidean-space) set
fix *x*::'a::euclidean-space
fix γ
assume *ind-hyp*: $(\forall b \in \text{basis}. \text{line-integral-exists } F \{b\} \gamma \implies \text{line-integral-exists } F \text{ basis } \gamma)$
 $(\forall b \in \text{basis}. \text{line-integral-exists } F \{b\} \gamma \implies \text{line-integral } F \text{ basis } \gamma = (\sum_{b \in \text{basis}} \text{line-integral } F \{b\} \gamma))$
assume *step*: *finite basis*
 $x \notin \text{basis}$
 $\forall b \in \text{insert } x \text{ basis}. \text{line-integral-exists } F \{b\} \gamma$
then have *0*: *line-integral-exists* $F \{x\} \gamma$ **by** *auto*
have *1*: *line-integral-exists* $F \text{ basis } \gamma$
using *ind-hyp step* **by** *auto*
then show *line-integral-exists* $F (\text{insert } x \text{ basis}) \gamma$
using *step(1) step(2) line-integral-sum-gen(2)*[*OF - 0 1*] **by** *simp*
have *3*: *finite (insert x basis)* **using** *step(1)* **by** *auto*
have *line-integral* $F \text{ basis } \gamma = (\sum_{b \in \text{basis}} \text{line-integral } F \{b\} \gamma)$
using *ind-hyp step* **by** *auto*
then show *line-integral* $F (\text{insert } x \text{ basis}) \gamma = (\sum_{b \in \text{insert } x \text{ basis}} \text{line-integral } F \{b\} \gamma)$
using *step(1) step(2) line-integral-sum-gen(1)*[*OF 3 0 1*]
by *force*
qed

lemma *reparam-weak-eq-line-integrals-basis*:

assumes *reparam-weak* $\gamma 1 \ \gamma 2$
 $\gamma 2$ *C1-differentiable-on* $\{0..1\}$
 $\forall b \in \text{basis}. \text{continuous-on } (\text{path-image } \gamma 2) (\lambda x. F x \cdot b)$
finite basis
shows *line-integral* $F \text{ basis } \gamma 1 = \text{line-integral } F \text{ basis } \gamma 2$
line-integral-exists $F \text{ basis } \gamma 1$
proof–
show *line-integral-exists* $F \text{ basis } \gamma 1$
using *reparam-weak-eq-line-integrals(2)*[*OF assms(1) assms(2)*] *assms(3-4)*
line-integral-sum-basis(2)[*OF assms(4)*]
by(*simp add: subset-iff*)
show *line-integral* $F \text{ basis } \gamma 1 = \text{line-integral } F \text{ basis } \gamma 2$
using *reparam-weak-eq-line-integrals*[*OF assms(1) assms(2)*] *assms(3-4)* *line-integral-sum-basis(1)*[*OF assms(4)*]
line-integral-exists-smooth-one-base[*OF assms(2)*]
by(*simp add: subset-iff*)
qed

lemma *reparam-eq-line-integrals-basis*:

assumes *reparam* $\gamma 1 \ \gamma 2$
 $\gamma 2$ *piecewise-C1-differentiable-on* $\{0..1\}$
 $\forall b \in \text{basis}. \text{continuous-on } (\text{path-image } \gamma 2) (\lambda x. F x \cdot b)$
finite basis

$\forall b \in \text{basis}. \text{line-integral-exists } F \{b\} \gamma 2$
shows $\text{line-integral } F \text{ basis } \gamma 1 = \text{line-integral } F \text{ basis } \gamma 2$
 $\text{line-integral-exists } F \text{ basis } \gamma 1$
proof –
show $\text{line-integral-exists } F \text{ basis } \gamma 1$
using $\text{reparam-eq-line-integrals}(2)[OF \text{ assms}(1) \text{ assms}(2)] \text{ assms}(3-5) \text{ line-integral-sum-basis}(2)[OF \text{ assms}(4)]$
by (*simp add: subset-iff*)
show $\text{line-integral } F \text{ basis } \gamma 1 = \text{line-integral } F \text{ basis } \gamma 2$
using $\text{reparam-eq-line-integrals}[OF \text{ assms}(1) \text{ assms}(2)] \text{ assms}(3-5) \text{ line-integral-sum-basis}(1)[OF \text{ assms}(4)]$
by (*simp add: subset-iff*)
qed

lemma *line-integral-exists-smooth*:
assumes $\gamma \text{ C1-differentiable-on } \{0..1\}$
 $\forall (b::'a::\text{euclidean-space}) \in \text{basis}. \text{continuous-on } (\text{path-image } \gamma) (\lambda x. F x \cdot b)$
 finite basis
shows $\text{line-integral-exists } F \text{ basis } \gamma$
using $\text{reparam-weak-eq-line-integrals-basis}(2)[OF \text{ reparam-weak-eq-refl}[where ?\gamma 1 . 0 = \gamma]] \text{ assms}$
by *fastforce*

lemma *smooth-path-imp-reverse*:
assumes $g \text{ C1-differentiable-on } \{0..1\}$
shows $(\text{reversepath } g) \text{ C1-differentiable-on } \{0..1\}$
using $\text{assms } \text{continuous-on-const}$
apply (*auto simp: reversepath-def*)
apply (*rule C1-differentiable-compose [of \lambda x::real. 1-x - g, unfolded o-def]*)
apply (*auto simp: C1-differentiable-on-eq*)
apply (*simp add: finite-vimageI inj-on-def*)
done

lemma *piecewise-smooth-path-imp-reverse*:
assumes $g \text{ piecewise-C1-differentiable-on } \{0..1\}$
shows $(\text{reversepath } g) \text{ piecewise-C1-differentiable-on } \{0..1\}$
using $\text{assms } \text{valid-path-reversepath}$
using *valid-path-def* **by** *blast*

definition *chain-reparam-weak-chain* **where**
 $\text{chain-reparam-weak-chain } \text{one-chain1 } \text{one-chain2} \equiv$
 $\exists f. \text{bij } f \wedge f ' \text{one-chain1} = \text{one-chain2} \wedge (\forall (k,\gamma) \in \text{one-chain1}. \text{if } k = \text{fst} (f(k,\gamma)) \text{ then } \text{reparam-weak } \gamma (\text{snd } (f(k,\gamma))) \text{ else } \text{reparam-weak } \gamma (\text{reversepath } (\text{snd } (f(k,\gamma)))))$

lemma *chain-reparam-weak-chain-line-integral*:
assumes $\text{chain-reparam-weak-chain } \text{one-chain1 } \text{one-chain2}$
 $\forall (k2,\gamma2) \in \text{one-chain2}. \gamma 2 \text{ C1-differentiable-on } \{0..1\}$
 $\forall (k2,\gamma2) \in \text{one-chain2}. \forall b \in \text{basis}. \text{continuous-on } (\text{path-image } \gamma 2) (\lambda x. F x \cdot b)$

finite basis
and *bound1: boundary-chain one-chain1*
and *bound2: boundary-chain one-chain2*
shows *one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis one-chain2*
 $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$
proof –
obtain *f where f: bij f*
 $(\forall (k, \gamma) \in \text{one-chain1}. \text{if } k = \text{fst } (f(k, \gamma)) \text{ then reparam-weak } \gamma (\text{snd } (f(k, \gamma))) \text{ else reparam-weak } \gamma (\text{reversepath } (\text{snd } (f(k, \gamma))))))$
 $f \text{ ' one-chain1 = one-chain2}$
using *assms(1)*
by *(auto simp add: chain-reparam-weak-chain-def)*
have $0: \forall x \in \text{one-chain1}. (\text{case } x \text{ of } (k, \gamma) \Rightarrow (\text{real-of-int } k * \text{line-integral } F \text{ basis } \gamma) = (\text{case } f \text{ x of } (k, \gamma) \Rightarrow \text{real-of-int } k * \text{line-integral } F \text{ basis } \gamma) \wedge \text{line-integral-exists } F \text{ basis } \gamma)$
proof
{fix *k1 γ 1*
assume *ass1: (k1, γ 1) \in one-chain1*
have $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma1 = (\text{case } (f (k1, \gamma1)) \text{ of } (k2, \gamma2) \Rightarrow \text{real-of-int } k2 * \text{line-integral } F \text{ basis } \gamma2) \wedge \text{line-integral-exists } F \text{ basis } \gamma1$
proof(*cases*)
assume *ass2: k1 = 1*
let *?k2 = fst (f (k1, γ 1))*
let *? γ 2 = snd (f (k1, γ 1))*
have $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma1 = \text{real-of-int } ?k2 * \text{line-integral } F \text{ basis } ?\gamma2 \wedge \text{line-integral-exists } F \text{ basis } \gamma1$
proof(*cases*)
assume *ass3: ?k2 = 1*
then have $0: \text{reparam-weak } \gamma1 ?\gamma2$
using *ass1 ass2 f(2)*
by *auto*
have $1: ?\gamma2 \text{ C1-differentiable-on } \{0..1\}$
using *f(3) assms(2) ass1*
by *force*
have $2: \forall b \in \text{basis}. \text{continuous-on } (\text{path-image } ?\gamma2) (\lambda x. F x \cdot b)$
using *f(3) assms(3) ass1*
by *force*
show $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma1 = \text{real-of-int } ?k2 * \text{line-integral } F \text{ basis } ?\gamma2 \wedge \text{line-integral-exists } F \text{ basis } \gamma1$
using *assms reparam-weak-eq-line-integrals-basis[OF 0 1 2 assms(4)]*
ass2 ass3
by *auto*
next
assume *?k2 \neq 1*
then have *ass3: ?k2 = -1*

```

    using bound2 ass1 f(3) unfolding boundary-chain-def by force
  then have 0: reparam-weak  $\gamma_1$  (reversepath  $?\gamma_2$ )
    using ass1 ass2 f(2)
    by auto
  have 1: (reversepath  $?\gamma_2$ ) C1-differentiable-on {0..1}
    using f(3) assms(2) ass1 smooth-path-imp-reverse
    by force
  have 2:  $\forall b \in \text{basis}. \text{continuous-on} (\text{path-image} (\text{reversepath } ?\gamma_2)) (\lambda x. F x$ 
· b)
    using f(3) assms(3) ass1 path-image-reversepath
    by force
  have 3: line-integral F basis  $?\gamma_2 = - \text{line-integral F basis} (\text{reversepath}
?\gamma_2)$ 
  proof -
    have i: valid-path (reversepath  $?\gamma_2$ )
      using 1 C1-differentiable-imp-piecewise
      by (auto simp add: valid-path-def)
    show ?thesis
      using line-integral-on-reverse-path(1)[OF i line-integral-exists-smooth[OF
1 2 ]] assms
      by auto
    qed
  show real-of-int  $k_1 * \text{line-integral F basis } \gamma_1 = \text{real-of-int } ?k_2 * \text{line-integral}
F \text{ basis } ?\gamma_2 \wedge$ 
      line-integral-exists F basis  $\gamma_1$ 
      using assms reparam-weak-eq-line-integrals-basis[OF 0 1 2 assms(4)]
      ass2 ass3 3
      by auto
    qed
  then show real-of-int  $k_1 * \text{line-integral F basis } \gamma_1 = (\text{case } f (k_1, \gamma_1) \text{ of}
(k_2, \gamma_2) \Rightarrow \text{real-of-int } k_2 * \text{line-integral F basis } \gamma_2) \wedge$ 
      line-integral-exists F basis  $\gamma_1$ 
      by (simp add: case-prod-beta')
  next
  assume  $k_1 \neq 1$ 
  then have ass2:  $k_1 = -1$ 
    using bound1 ass1 f(3) unfolding boundary-chain-def by force
  let  $?k_2 = \text{fst } (f (k_1, \gamma_1))$ 
  let  $?\gamma_2 = \text{snd } (f (k_1, \gamma_1))$ 
  have real-of-int  $k_1 * \text{line-integral F basis } \gamma_1 = \text{real-of-int } ?k_2 * \text{line-integral}
F \text{ basis } ?\gamma_2 \wedge$ 
      line-integral-exists F basis  $\gamma_1$ 
  proof(cases)
    assume ass3:  $?k_2 = 1$ 
    then have 0: reparam-weak  $\gamma_1$  (reversepath  $?\gamma_2$ )
      using ass1 ass2 f(2)
      by auto
    have 1: (reversepath  $?\gamma_2$ ) C1-differentiable-on {0..1}
      using f(3) assms(2) ass1 smooth-path-imp-reverse

```

```

    by force
  have 2:  $\forall b \in \text{basis}. \text{continuous-on } (\text{path-image } (\text{reversepath } ?\gamma 2)) (\lambda x. F x$ 
  · b)
    using f(3) assms(3) ass1 path-image-reversepath
    by force
  have 3:  $\text{line-integral } F \text{ basis } ?\gamma 2 = - \text{line-integral } F \text{ basis } (\text{reversepath}$ 
   $?\gamma 2)$ 
  proof -
    have i:  $\text{valid-path } (\text{reversepath } ?\gamma 2)$ 
      using 1 C1-differentiable-imp-piecewise
      by (auto simp add: valid-path-def)
    show ?thesis
      using line-integral-on-reverse-path(1)[OF i line-integral-exists-smooth[OF
  1 2 assms(4)]]
      by auto
    qed
    show  $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma 1 = \text{real-of-int } ?k2 * \text{line-integral}$ 
   $F \text{ basis } ?\gamma 2 \wedge$ 
       $\text{line-integral-exists } F \text{ basis } \gamma 1$ 
      using assms reparam-weak-eq-line-integrals-basis[OF 0 1 2 assms(4)]
      ass2 ass3 3
      by auto
  next
    assume  $?k2 \neq 1$ 
    then have  $\text{ass3}: ?k2 = -1$ 
      using bound2 ass1 f(3) unfolding boundary-chain-def by force
    then have 0:  $\text{reparam-weak } \gamma 1 ?\gamma 2$ 
      using ass1 ass2 f(2) by auto
    have 1:  $?\gamma 2 \text{ C1-differentiable-on } \{0..1\}$ 
      using f(3) assms(2) ass1 by force
    have 2:  $\forall b \in \text{basis}. \text{continuous-on } (\text{path-image } ?\gamma 2) (\lambda x. F x \cdot b)$ 
      using f(3) assms(3) ass1 by force
    show  $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma 1 = \text{real-of-int } ?k2 * \text{line-integral}$ 
   $F \text{ basis } ?\gamma 2 \wedge$ 
       $\text{line-integral-exists } F \text{ basis } \gamma 1$ 
      using assms reparam-weak-eq-line-integrals-basis[OF 0 1 2 assms(4)]
      ass2 ass3
      by auto
    qed
    then show  $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma 1 = (\text{case } f (k1, \gamma 1) \text{ of}$ 
   $(k2, \gamma 2) \Rightarrow \text{real-of-int } k2 * \text{line-integral } F \text{ basis } \gamma 2) \wedge$ 
       $\text{line-integral-exists } F \text{ basis } \gamma 1$ 
      by (simp add: case-prod-beta')
    qed
  }
  then show  $\bigwedge x. x \in \text{one-chain1} \Longrightarrow$ 
       $(\text{case } x \text{ of } (k, \gamma) \Rightarrow (\text{real-of-int } k * \text{line-integral } F \text{ basis}$ 
   $\gamma) = (\text{case } f x \text{ of } (k, \gamma) \Rightarrow \text{real-of-int } k * \text{line-integral } F \text{ basis } \gamma) \wedge$ 
       $\text{line-integral-exists } F \text{ basis } \gamma)$ 

```


by (auto simp add: case-prod-beta')
 qed
 show one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis
 one-chain2
 using 0 by (simp add: one-chain-line-integral-def sum-bij[OF f(1) - f(3)]
 split-beta)
 show $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$
 using 0 by blast
 qed

definition chain-reparam-chain where

chain-reparam-chain one-chain1 one-chain2 \equiv
 $\exists f. \text{bij } f \wedge f' \text{ one-chain1} = \text{one-chain2} \wedge (\forall (k, \gamma) \in \text{one-chain1}. \text{if } k = \text{fst}$
 $(f(k, \gamma)) \text{ then reparam } \gamma (\text{snd } (f(k, \gamma))) \text{ else reparam } \gamma (\text{reversepath } (\text{snd } (f(k, \gamma))))))$

definition chain-reparam-weak-path::((real) \Rightarrow (real * real)) \Rightarrow ((int * ((real) \Rightarrow
 (real * real))) set) \Rightarrow bool where

chain-reparam-weak-path γ one-chain
 $\equiv \exists l. \text{set } l = \text{one-chain} \wedge \text{distinct } l \wedge \text{reparam } \gamma (\text{rec-join } l) \wedge \text{valid-chain-list}$
 $l \wedge l \neq []$

lemma chain-reparam-chain-line-integral:

assumes chain-reparam-chain one-chain1 one-chain2
 $\forall (k2, \gamma2) \in \text{one-chain2}. \gamma2 \text{ piecewise-C1-differentiable-on } \{0..1\}$
 $\forall (k2, \gamma2) \in \text{one-chain2}. \forall b \in \text{basis}. \text{continuous-on } (\text{path-image } \gamma2) (\lambda x. F x \cdot b)$
 finite basis
 and bound1: boundary-chain one-chain1
 and bound2: boundary-chain one-chain2
 and line: $\forall (k2, \gamma2) \in \text{one-chain2}. (\forall b \in \text{basis}. \text{line-integral-exists } F \{b\} \gamma2)$
 shows one-chain-line-integral F basis one-chain1 = one-chain-line-integral F
 basis one-chain2
 $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$

proof–

obtain f where f: bij f
 $(\forall (k, \gamma) \in \text{one-chain1}. \text{if } k = \text{fst } (f(k, \gamma)) \text{ then reparam } \gamma (\text{snd } (f(k, \gamma))) \text{ else}$
 $\text{reparam } \gamma (\text{reversepath } (\text{snd } (f(k, \gamma))))))$
 $f' \text{ one-chain1} = \text{one-chain2}$
 using assms(1)
 by (auto simp add: chain-reparam-chain-def)
 have integ-exist-b: $\forall (k1, \gamma1) \in \text{one-chain1}. \forall b \in \text{basis}. \text{line-integral-exists } F \{b\}$
 $(\text{snd } (f(k1, \gamma1)))$
 using line f by fastforce
 have valid-cubes: $\forall (k1, \gamma1) \in \text{one-chain1}. \text{valid-path } (\text{snd } (f(k1, \gamma1)))$
 using assms(2) f(3) valid-path-def by fastforce
 have integ-rev-exist-b: $\forall (k1, \gamma1) \in \text{one-chain1}. \forall b \in \text{basis}. \text{line-integral-exists } F \{b\}$
 $(\text{reversepath } (\text{snd } (f(k1, \gamma1))))$
 using line-integral-on-reverse-path(2) integ-exist-b valid-cubes
 by blast
 have 0: $\forall x \in \text{one-chain1}. (\text{case } x \text{ of } (k, \gamma) \Rightarrow (\text{real-of-int } k * \text{line-integral } F \text{ basis}$

$\gamma) = (\text{case } f \text{ of } (k, \gamma) \Rightarrow \text{real-of-int } k * \text{line-integral } F \text{ basis } \gamma) \wedge$
 $\text{line-integral-exists } F \text{ basis } \gamma)$

proof
 {fix $k1 \ \gamma1$
 assume $ass1: (k1, \gamma1) \in \text{one-chain1}$
 have $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma1 = (\text{case } (f \ (k1, \gamma1)) \text{ of } (k2, \gamma2)$
 $\Rightarrow \text{real-of-int } k2 * \text{line-integral } F \text{ basis } \gamma2) \wedge$
 $\text{line-integral-exists } F \text{ basis } \gamma1$

proof(cases)
 assume $ass2: k1 = 1$
 let $?k2 = \text{fst } (f \ (k1, \gamma1))$
 let $? \gamma2 = \text{snd } (f \ (k1, \gamma1))$
 have $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma1 = \text{real-of-int } ?k2 * \text{line-integral}$
 $F \text{ basis } ? \gamma2 \wedge$
 $\text{line-integral-exists } F \text{ basis } \gamma1$

proof(cases)
 assume $ass3: ?k2 = 1$
 then have $0: \text{reparam } \gamma1 \ ? \gamma2$
 using $ass1 \ ass2 \ f(2)$
 by auto
 have $1: ? \gamma2 \text{ piecewise-}C1\text{-differentiable-on } \{0..1\}$
 using $f(3) \ assms(2) \ ass1$
 by force
 have $2: \forall b \in \text{basis. continuous-on } (\text{path-image } ? \gamma2) \ (\lambda x. F \ x \cdot b)$
 using $f(3) \ assms(3) \ ass1$
 by force
 show $\text{real-of-int } k1 * \text{line-integral } F \text{ basis } \gamma1 = \text{real-of-int } ?k2 * \text{line-integral}$
 $F \text{ basis } ? \gamma2 \wedge$
 $\text{line-integral-exists } F \text{ basis } \gamma1$

using $assms \ \text{reparam-eq-line-integrals-basis}[OF \ 0 \ 1 \ 2 \ assms(4)] \ \text{integ-exist-b}$
 $ass1 \ ass2 \ ass3$
 by auto

next
 assume $?k2 \neq 1$
 then have $ass3: ?k2 = -1$
 using $\text{bound2 } ass1 \ f(3) \ \text{unfolding } \text{boundary-chain-def} \ \text{by force}$
 then have $0: \text{reparam } \gamma1 \ (\text{reversepath } ? \gamma2)$
 using $ass1 \ ass2 \ f(2)$
 by auto
 have $1: (\text{reversepath } ? \gamma2) \text{ piecewise-}C1\text{-differentiable-on } \{0..1\}$
 using $f(3) \ assms(2) \ ass1 \ \text{piecewise-smooth-path-imp-reverse}$
 by force
 have $2: \forall b \in \text{basis. continuous-on } (\text{path-image } (\text{reversepath } ? \gamma2)) \ (\lambda x. F \ x$
 $\cdot b)$
 using $f(3) \ assms(3) \ ass1 \ \text{path-image-reversepath}$
 by force
 have $3: \text{line-integral } F \text{ basis } ? \gamma2 = - \text{line-integral } F \text{ basis } (\text{reversepath}$
 $? \gamma2)$

proof–

```

have i: valid-path (reversepath ? $\gamma$ 2)
  using 1 C1-differentiable-imp-piecewise
  by (auto simp add: valid-path-def)
have ii: line-integral-exists F basis (snd (f (k1,  $\gamma$ 1)))
  using assms(4) line-integral-sum-basis(2) integ-exist-b ass1
  by fastforce
show ?thesis
  using i ii line-integral-on-reverse-path(1) valid-path-reversepath by blast
qed
show real-of-int k1 * line-integral F basis  $\gamma$ 1 = real-of-int ?k2 * line-integral
F basis ? $\gamma$ 2  $\wedge$ 
      line-integral-exists F basis  $\gamma$ 1
  using assms reparam-eq-line-integrals-basis[OF 0 1 2 assms(4)] in-
teg-rev-exist-b
      ass1 ass2 ass3 3
  by auto
qed
then show real-of-int k1 * line-integral F basis  $\gamma$ 1 = (case f (k1,  $\gamma$ 1) of
(k2,  $\gamma$ 2)  $\Rightarrow$  real-of-int k2 * line-integral F basis  $\gamma$ 2)  $\wedge$ 
      line-integral-exists F basis  $\gamma$ 1
  by (simp add: case-prod-beta')
next
assume k1  $\neq$  1
then have ass2: k1 = -1
  using bound1 ass1 f(3) unfolding boundary-chain-def by force
let ?k2 = fst (f (k1,  $\gamma$ 1))
let ? $\gamma$ 2 = snd (f (k1,  $\gamma$ 1))
have real-of-int k1 * line-integral F basis  $\gamma$ 1 = real-of-int ?k2 * line-integral
F basis ? $\gamma$ 2  $\wedge$ 
      line-integral-exists F basis  $\gamma$ 1
proof(cases)
assume ass3: ?k2 = 1
then have 0: reparam  $\gamma$ 1 (reversepath ? $\gamma$ 2)
  using ass1 ass2 f(2)
  by auto
have 1: (reversepath ? $\gamma$ 2) piecewise-C1-differentiable-on {0..1}
  using f(3) assms(2) ass1 piecewise-smooth-path-imp-reverse
  by force
have 2:  $\forall b \in \text{basis. continuous-on (path-image (reversepath ? $\gamma$ 2)) (\lambda x. F x$ 
. b)
  using f(3) assms(3) ass1 path-image-reversepath
  by force
have 3: line-integral F basis ? $\gamma$ 2 = - line-integral F basis (reversepath
? $\gamma$ 2)
proof-
have i: valid-path (reversepath ? $\gamma$ 2)
  using 1 C1-differentiable-imp-piecewise
  by (auto simp add: valid-path-def)
show ?thesis

```

```

    using line-integral-on-reverse-path(1)[OF i] integ-rev-exist-b
    using ass1 assms(4) line-integral-sum-basis(2) by fastforce
  qed
  show real-of-int k1 * line-integral F basis  $\gamma_1$  = real-of-int ?k2 * line-integral
  F basis ? $\gamma_2$   $\wedge$ 
    line-integral-exists F basis  $\gamma_1$ 
    using assms reparam-eq-line-integrals-basis[OF 0 1 2 assms(4)]
    ass2 ass3 3
    using ass1 integ-rev-exist-b by auto
  next
  assume ?k2  $\neq$  1
  then have ass3: ?k2 = -1
    using bound2 ass1 f(3) unfolding boundary-chain-def by force
  then have 0: reparam  $\gamma_1$  ? $\gamma_2$ 
    using ass1 ass2 f(2) by auto
  have 1: ? $\gamma_2$  piecewise-C1-differentiable-on {0..1}
    using f(3) assms(2) ass1
    by force
  have 2:  $\forall b \in \text{basis}. \text{continuous-on } (\text{path-image } ?\gamma_2) (\lambda x. F x \cdot b)$ 
    using f(3) assms(3) ass1
    by force
  show real-of-int k1 * line-integral F basis  $\gamma_1$  = real-of-int ?k2 * line-integral
  F basis ? $\gamma_2$   $\wedge$ 
    line-integral-exists F basis  $\gamma_1$ 
    using assms reparam-eq-line-integrals-basis[OF 0 1 2 assms(4)]
    ass2 ass3
    using ass1 integ-exist-b by auto
  qed
  then show real-of-int k1 * line-integral F basis  $\gamma_1$  = (case f (k1,  $\gamma_1$ ) of
  (k2,  $\gamma_2$ )  $\Rightarrow$  real-of-int k2 * line-integral F basis  $\gamma_2$ )  $\wedge$ 
    line-integral-exists F basis  $\gamma_1$ 
    by (simp add: case-prod-beta')
  qed
}
then show  $\bigwedge x. x \in \text{one-chain1} \Rightarrow$ 
  (case x of (k,  $\gamma$ )  $\Rightarrow$  (real-of-int k * line-integral F basis
 $\gamma$ ) = (case f x of (k,  $\gamma$ )  $\Rightarrow$  real-of-int k * line-integral F basis  $\gamma$ )  $\wedge$ 
  line-integral-exists F basis  $\gamma$ )
  by (auto simp add: case-prod-beta')
qed
show one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis
one-chain2
  using 0 by (simp add: one-chain-line-integral-def sum-bij[OF f(1) - f(3)]
  prod.case-eq-if)
  show  $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$ 
  using 0 by blast
qed
lemma path-image-rec-join:

```

```

fixes  $\gamma::real \Rightarrow (real \times real)$ 
fixes  $k::int$ 
fixes  $l$ 
shows  $\bigwedge k \gamma. (k, \gamma) \in set\ l \Longrightarrow valid-chain-list\ l \Longrightarrow path-image\ \gamma \subseteq path-image$ 
 $(rec-join\ l)$ 
proof(induction l)
  case Nil
  then show ?case by auto
next
  case ass: (Cons a l)
  obtain  $k' \gamma'$  where  $a = (k', \gamma')$  by force
  have  $path-image\ \gamma \subseteq path-image\ (rec-join\ ((k', \gamma') \# l))$ 
  proof(cases)
    assume  $l = []$ 
    then show  $path-image\ \gamma \subseteq path-image\ (rec-join\ ((k', \gamma') \# l))$ 
    using  $ass(2-3)\ a$ 
    by(auto simp add:)
  next
    assume  $l \neq []$ 
    then obtain  $b\ l'$  where  $b-l: l = b \# l'$ 
    by (meson rec-join.elims)
    obtain  $k'' \gamma''$  where  $b = (k'', \gamma'')$  by force
    show ?thesis
    using  $ass\ path-image-reversepath\ b-l\ path-image-join$ 
    by(fastforce simp add: a)
  qed
  then show ?case
  using  $a$  by auto
qed

```

lemma *path-image-rec-join-2:*

```

fixes  $l$ 
shows  $l \neq [] \Longrightarrow valid-chain-list\ l \Longrightarrow path-image\ (rec-join\ l) \subseteq (\bigcup (k, \gamma) \in set$ 
 $l. path-image\ \gamma)$ 
proof(induction l)
  case Nil
  then show ?case by auto
next
  case ass: (Cons a l)
  obtain  $k' \gamma'$  where  $a = (k', \gamma')$  by force
  have  $path-image\ (rec-join\ (a \# l)) \subseteq (\bigcup (k, y) \in set\ (a \# l). path-image\ y)$ 
  proof(cases)
    assume  $l = []$ 
    then show  $path-image\ (rec-join\ (a \# l)) \subseteq (\bigcup (k, y) \in set\ (a \# l). path-image$ 
 $y)$ 
    using  $step\ a$  by(auto simp add:)
  next
    assume  $l \neq []$ 
    then obtain  $b\ l'$  where  $b-l: l = b \# l'$ 

```

```

    by (meson rec-join.elims)
  obtain  $k'' \gamma''$  where  $b: b = (k'', \gamma'')$  by force
  show ?thesis
    using ass
      path-image-reversepath b-l path-image-join
    apply (auto simp add: a)
    apply blast
    by fastforce
  qed
  then show ?case
    using a by auto
  qed

```

```

lemma continuous-on-closed-UN:
  assumes finite S
  shows (( $\bigwedge s. s \in S \implies \text{closed } s$ )  $\implies$  ( $\bigwedge s. s \in S \implies \text{continuous-on } s f$ )  $\implies$ 
  continuous-on ( $\bigcup S$ ) f)
  using assms
  proof (induction S)
    case empty
    then show ?case by auto
  next
    case (insert x F)
    then show ?case using continuous-on-closed-Un closed-Union
      by (simp add: closed-Union continuous-on-closed-Un)
  qed

```

```

lemma chain-reparam-weak-path-line-integral:
  assumes path-eq-chain: chain-reparam-weak-path  $\gamma$  one-chain and
    boundary-chain: boundary-chain one-chain and
    line-integral-exists:  $\forall b \in \text{basis}. \forall (k :: \text{int}, \gamma) \in \text{one-chain}. \text{line-integral-exists } F \{b\}$ 
   $\gamma$  and
    valid-path:  $\forall (k :: \text{int}, \gamma) \in \text{one-chain}. \text{valid-path } \gamma$  and
    finite-basis: finite basis and
    cont:  $\forall b \in \text{basis}. \forall (k, \gamma 2) \in \text{one-chain}. \text{continuous-on } (\text{path-image } \gamma 2) (\lambda x. F x \cdot$ 
  b) and
    finite-one-chain: finite one-chain
  shows line-integral F basis  $\gamma = \text{one-chain-line-integral } F \text{ basis one-chain}$ 
    line-integral-exists F basis  $\gamma$ 

```

```

proof -
  obtain l where l-props: set l = one-chain distinct l reparam  $\gamma$  (rec-join l)
  valid-chain-list l l  $\neq []$ 
  using chain-reparam-weak-path-def assms
  by auto
  have bchain-l: boundary-chain (set l)
  using l-props boundary-chain
  by (simp add: boundary-chain-def)
  have cont-forall:  $\forall b \in \text{basis}. \text{continuous-on } (\bigcup (k, \gamma) \in \text{one-chain}. \text{path-image } \gamma)$ 

```

```

( $\lambda x. F x \cdot b$ )
proof
  fix  $b$ 
  assume  $ass: b \in basis$ 
  have  $continuous-on (\bigcup ((path-image \circ snd) \text{ `one-chain})) (\lambda x. F x \cdot b)$ 
    apply( $rule\ continuous-on-closed-UN[\text{where } ?S = (path-image \circ snd) \text{ `one-chain}]$ )
  proof
    show  $finite\ one-chain$  using  $finite-one-chain$  by  $auto$ 
    show  $\bigwedge s. s \in (path-image \circ snd) \text{ `one-chain} \implies closed\ s$ 
      using  $closed-valid-path-image[OF]\ valid-path$ 
      by  $fastforce$ 
    show  $\bigwedge s. s \in (path-image \circ snd) \text{ `one-chain} \implies continuous-on\ s (\lambda x. F x \cdot b)$ 
  using  $cont\ ass$  by  $force$ 
qed
then show  $continuous-on (\bigcup (k, \gamma) \in one-chain. path-image\ \gamma) (\lambda x. F x \cdot b)$ 
  by ( $auto\ simp\ add: Union-eq\ case-prod-beta$ )
qed
show  $line-integral-exists\ F\ basis\ \gamma$ 
proof ( $rule\ reparam-eq-line-integrals-basis[OF\ l-props(3) - -\ finite-basis]$ )
  let  $? \gamma 2.0 = rec-join\ l$ 
  show  $? \gamma 2.0\ piecewise-C1-differentiable-on\ \{0..1\}$ 
    apply( $simp\ only: valid-path-def[symmetric]$ )
    apply( $rule\ joined-is-valid$ )
    using  $assms\ l-props$  by  $auto$ 
  show  $\forall b \in basis. continuous-on (path-image (rec-join\ l)) (\lambda x. F x \cdot b)$  using
 $path-image-rec-join-2[OF\ l-props(5)\ l-props(4)]\ l-props(1)$ 
    using  $cont-forall\ continuous-on-subset$  by  $blast$ 
  show  $\forall b \in basis. line-integral-exists\ F\ \{b\} (rec-join\ l)$ 
proof
  fix  $b$ 
  assume  $ass: b \in basis$ 
  show  $line-integral-exists\ F\ \{b\} (rec-join\ l)$ 
proof ( $rule\ line-integral-exists-on-rec-join$ )
  show  $boundary-chain (set\ l)$ 
    using  $l-props\ boundary-chain$  by  $auto$ 
  show  $valid-chain-list\ l$  using  $l-props$  by  $auto$ 
  show  $\bigwedge k\ \gamma. (k, \gamma) \in set\ l \implies valid-path\ \gamma$  using  $l-props\ assms$  by  $auto$ 
  show  $\forall (k, \gamma) \in set\ l. line-integral-exists\ F\ \{b\}\ \gamma$  using  $l-props\ line-integral-exists$ 
 $ass$  by  $blast$ 
qed
qed
show  $line-integral\ F\ basis\ \gamma = one-chain-line-integral\ F\ basis\ one-chain$ 
proof -
  have  $i: line-integral\ F\ basis (rec-join\ l) = one-chain-line-integral\ F\ basis\ one-chain$ 
proof ( $rule\ one-chain-line-integral-rec-join$ )

```

```

show set l = one-chain distinct l valid-chain-list l using l-props by auto
show boundary-chain one-chain using boundary-chain by auto
show  $\forall (k, \gamma) \in \text{one-chain}. \text{line-integral-exists } F \text{ basis } \gamma$ 
  using line-integral-sum-basis(2)[OF finite-basis] line-integral-exists by blast
show  $\forall (k, \gamma) \in \text{one-chain}. \text{valid-path } \gamma$  using valid-path by auto
show finite basis using finite-basis by auto
qed
have ii: line-integral F basis  $\gamma = \text{line-integral } F \text{ basis } (\text{rec-join } l)$ 
proof (rule reparam-eq-line-integrals-basis[OF l-props(3) - - finite-basis])
  let ? $\gamma$ 2.0=rec-join l
  show ? $\gamma$ 2.0 piecewise-C1-differentiable-on {0..1}
    apply(simp only: valid-path-def[symmetric])
    apply(rule joined-is-valid)
    using assms l-props by auto
  show  $\forall b \in \text{basis}. \text{continuous-on } (\text{path-image } (\text{rec-join } l)) (\lambda x. F x \cdot b)$  using
path-image-rec-join-2[OF l-props(5) l-props(4)] l-props(1)
    using cont-forall continuous-on-subset by blast
  show  $\forall b \in \text{basis}. \text{line-integral-exists } F \{b\} (\text{rec-join } l)$ 
proof
  fix b
  assume ass:  $b \in \text{basis}$ 
  show line-integral-exists F {b} (rec-join l)
  proof (rule line-integral-exists-on-rec-join)
    show boundary-chain (set l)
      using l-props boundary-chain by auto
    show valid-chain-list l using l-props by auto
    show  $\bigwedge k \gamma. (k, \gamma) \in \text{set } l \implies \text{valid-path } \gamma$  using l-props assms by auto
    show  $\forall (k, \gamma) \in \text{set } l. \text{line-integral-exists } F \{b\} \gamma$  using l-props line-integral-exists
ass by blast
  qed
qed
qed
show line-integral F basis  $\gamma = \text{one-chain-line-integral } F \text{ basis one-chain}$  using
i ii by auto
qed
qed

```

definition chain-reparam-chain' **where**
 chain-reparam-chain' one-chain1 subdiv
 $\equiv \exists f. ((\bigcup (f \text{ ' one-chain1})) = \text{subdiv}) \wedge$
 $(\forall \text{cube} \in \text{one-chain1}. \text{chain-reparam-weak-path } (\text{rec-join } [\text{cube}]) (f \text{ cube}))$
 \wedge
 $(\forall p \in \text{one-chain1}. \forall p' \in \text{one-chain1}. p \neq p' \implies f p \cap f p' = \{\}) \wedge$
 $(\forall x \in \text{one-chain1}. \text{finite } (f x))$

lemma chain-reparam-chain'-imp-finite-subdiv:
assumes finite one-chain1
 chain-reparam-chain' one-chain1 subdiv
shows finite subdiv


```

using assms
by(auto simp add: chain-reparam-chain'-def)

lemma chain-reparam-chain'-line-integral:
  assumes chain1-eq-chain2: chain-reparam-chain' one-chain1 subdiv and
    boundary-chain1: boundary-chain one-chain1 and
    boundary-chain2: boundary-chain subdiv and
    line-integral-exists-on-chain2:  $\forall b \in \text{basis}. \forall (k::\text{int}, \gamma) \in \text{subdiv}. \text{line-integral-exists}$ 
F {b}  $\gamma$  and
    valid-path:  $\forall (k, \gamma) \in \text{subdiv}. \text{valid-path } \gamma$  and
    valid-path-2:  $\forall (k, \gamma) \in \text{one-chain1}. \text{valid-path } \gamma$  and
    finite-chain1: finite one-chain1 and
    finite-basis: finite basis and
    cont-field:  $\forall b \in \text{basis}. \forall (k, \gamma 2) \in \text{subdiv}. \text{continuous-on } (\text{path-image } \gamma 2) (\lambda x. F$ 
x  $\cdot$  b)
  shows one-chain-line-integral F basis one-chain1 = one-chain-line-integral F
basis subdiv
     $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$ 
proof –
  obtain f where f-props:
    ( $\bigcup (f \text{ ' one-chain1}) = \text{subdiv}$ )
    ( $\forall \text{cube} \in \text{one-chain1}. \text{chain-reparam-weak-path } (\text{rec-join } [\text{cube}]) (f \text{ cube})$ )
    ( $\forall p \in \text{one-chain1}. \forall p' \in \text{one-chain1}. p \neq p' \longrightarrow f p \cap f p' = \{\}$ )
    ( $\forall x \in \text{one-chain1}. \text{finite } (f x)$ )
  using chain1-eq-chain2
  by (auto simp add: chain-reparam-chain'-def)
  then have 0: one-chain-line-integral F basis subdiv = one-chain-line-integral F
basis ( $\bigcup (f \text{ ' one-chain1})$ )
  by auto
  {fix k  $\gamma$ 
  assume ass: (k,  $\gamma$ )  $\in$  one-chain1
  have line-integral-exists F basis  $\gamma \wedge$ 
    one-chain-line-integral F basis (f (k,  $\gamma$ )) = k * line-integral F basis  $\gamma$ 
proof(cases k = 1)
  assume k-eq-1: k = 1
  then have i: chain-reparam-weak-path  $\gamma$  (f(k, $\gamma$ ))
  using f-props(2) ass by auto
  have ii: boundary-chain (f(k, $\gamma$ ))
  using f-props(1) ass assms unfolding boundary-chain-def
  by blast
  have iii:  $\forall b \in \text{basis}. \forall (k, \gamma) \in f (k, \gamma). \text{line-integral-exists } F \{b\} \gamma$ 
  using f-props(1) ass assms
  by blast
  have iv:  $\forall (k, \gamma) \in f (k, \gamma). \text{valid-path } \gamma$ 
  using f-props(1) ass assms
  by blast
  have v:  $\forall b \in \text{basis}. \forall (k, \gamma 2) \in f (k, \gamma). \text{continuous-on } (\text{path-image } \gamma 2) (\lambda x. F$ 
x  $\cdot$  b)
  using f-props(1) ass assms

```

```

    by blast
    have line-integral-exists F basis  $\gamma \wedge$  one-chain-line-integral F basis  $(f (k, \gamma))$ 
= line-integral F basis  $\gamma$ 
    using chain-reparam-weak-path-line-integral[OF i ii iii iv finite-basis v] ass
f-props(4)
    by (auto)
    then show line-integral-exists F basis  $\gamma \wedge$  one-chain-line-integral F basis  $(f$ 
 $(k, \gamma)) = k * \text{line-integral F basis } \gamma$ 
    using k-eq-1 by auto
next
assume  $k \neq 1$ 
then have k-eq-neg1:  $k = -1$ 
    using ass boundary-chain1
    by (auto simp add: boundary-chain-def)
then have i:chain-reparam-weak-path (reversepath  $\gamma$ )  $(f(k,\gamma))$ 
    using f-props(2) ass by auto
have ii:boundary-chain  $(f(k,\gamma))$ 
    using f-props(1) ass assms unfolding boundary-chain-def
    by blast
have iii: $\forall b \in \text{basis}. \forall (k, \gamma) \in f (k, \gamma). \text{line-integral-exists F } \{b\} \gamma$ 
    using f-props(1) ass assms
    by blast
have iv:  $\forall (k, \gamma) \in f (k, \gamma). \text{valid-path } \gamma$ 
    using f-props(1) ass assms by blast
have v:  $\forall b \in \text{basis}. \forall (k, \gamma) \in f (k, \gamma). \text{continuous-on (path-image } \gamma) (\lambda x. F$ 
 $x \cdot b)$ 
    using f-props(1) ass assms by blast
have x:line-integral-exists F basis (reversepath  $\gamma$ )  $\wedge$  one-chain-line-integral F
basis  $(f (k, \gamma)) = \text{line-integral F basis (reversepath } \gamma)$ 
    using chain-reparam-weak-path-line-integral[OF i ii iii iv finite-basis v] ass
f-props(4)
    by auto
have valid-path (reversepath  $\gamma$ )
    using f-props(1) ass assms by auto
then have line-integral-exists F basis  $\gamma \wedge$  one-chain-line-integral F basis  $(f$ 
 $(k, \gamma)) = - (\text{line-integral F basis } \gamma)$ 
    using line-integral-on-reverse-path reversepath-reversepath x ass
    by metis
then show line-integral-exists F basis  $\gamma \wedge$  one-chain-line-integral F basis  $(f$ 
 $(k::\text{int}, \gamma)) = k * \text{line-integral F basis } \gamma$ 
    using k-eq-neg1 by auto
qed}
note cube-line-integ = this
have finite-chain2: finite subdiv
    using finite-chain1 f-props(1) f-props(4) by auto
show line-integral-exists-on-chain1:  $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists F}$ 
basis  $\gamma$ 
    using cube-line-integ by auto
have 1: one-chain-line-integral F basis  $(\bigcup (f ' \text{one-chain1})) = \text{one-chain-line-integral}$ 

```

F basis one-chain1
proof –
have $0: \text{one-chain-line-integral } F \text{ basis } (\bigcup (f \text{ ' one-chain1})) =$
 $(\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral } F$
basis one-chain)
proof –
have *finite*: $\forall \text{chain} \in (f \text{ ' one-chain1}). \text{finite chain}$
using *f-props(1) finite-chain2*
by (*meson Sup-upper finite-subset*)
have *disj*: $\forall A \in f \text{ ' one-chain1}. \forall B \in f \text{ ' one-chain1}. A \neq B \longrightarrow A \cap B = \{\}$
apply (*simp add:image-def*)
using *f-props(3)*
by *metis*
show $\text{one-chain-line-integral } F \text{ basis } (\bigcup (f \text{ ' one-chain1})) =$
 $(\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral}$
F basis one-chain)
using *Groups-Big.comm-monoid-add-class.sum.Union-disjoint[OF finite disj]*
one-chain-line-integral-def
by *auto*
qed
have $1: (\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral } F \text{ basis one-chain})$
 $=$
 $\text{one-chain-line-integral } F \text{ basis one-chain1}$
proof –
have $(\sum \text{one-chain} \in (f \text{ ' one-chain1}). \text{one-chain-line-integral } F \text{ basis one-chain})$
 $=$
 $(\sum (k, \gamma) \in \text{one-chain1}. k * (\text{line-integral } F \text{ basis } \gamma))$
proof –
have $i: (\sum \text{one-chain} \in (f \text{ ' } (\text{one-chain1} - \{p. f p = \{\}\})). \text{one-chain-line-integral}$
F basis one-chain) =
 $(\sum (k, \gamma) \in \text{one-chain1} - \{p. f p = \{\}\}. k * (\text{line-integral}$
*F basis } \gamma))
proof –
have *inj-on f* $(\text{one-chain1} - \{p. f p = \{\}\})$
unfolding *inj-on-def*
using *f-props(3) by force*
then have $0: (\sum \text{one-chain} \in (f \text{ ' } (\text{one-chain1} - \{p. f p = \{\}\})).$
one-chain-line-integral } F \text{ basis one-chain)
 $= (\sum (k, \gamma) \in (\text{one-chain1} - \{p. f p =$
 $\{\}\}). \text{one-chain-line-integral } F \text{ basis } (f (k, \gamma)))$
using *Groups-Big.comm-monoid-add-class.sum.reindex*
by *auto*
have $\bigwedge k \gamma. (k, \gamma) \in (\text{one-chain1} - \{p. f p = \{\}\}) \implies$
 $\text{one-chain-line-integral } F \text{ basis } (f(k, \gamma))$
 $= k * (\text{line-integral } F \text{ basis } \gamma)$
using *cube-line-integ by auto*
then have $(\sum (k, \gamma) \in (\text{one-chain1} - \{p. f p = \{\}\}). \text{one-chain-line-integral}$
F basis } (f (k, \gamma)))
 $= (\sum (k, \gamma) \in (\text{one-chain1} - \{p. f p = \{\}\}). k * (\text{line-integral } F$*

```

basis  $\gamma$ )
  by (auto intro!: Finite-Cartesian-Product.sum-cong-aux)
  then show  $(\sum one-chain \in (f' (one-chain1 - \{p. f p = \{\}\})). one-chain-line-integral$ 
 $F$  basis  $one-chain) =$ 
 $(\sum (k, \gamma) \in (one-chain1 - \{p. f p =$ 
 $\{\}\}). k * (line-integral F$  basis  $\gamma))$ 
    using 0 by auto
  qed
  have  $\bigwedge (k::int) \gamma. (k, \gamma) \in one-chain1 \implies (f (k, \gamma) = \{\}) \implies (k, \gamma) \in$ 
 $\{(k', \gamma'). k' * (line-integral F$  basis  $\gamma') = 0\}$ 
  proof-
    fix  $k::int$ 
    fix  $\gamma::real \implies real * real$ 
    assume  $ass:(k, \gamma) \in one-chain1$ 
     $(f (k, \gamma) = \{\})$ 
    then have  $zero-line-integral:one-chain-line-integral F$  basis  $(f (k, \gamma)) = 0$ 
      using  $one-chain-line-integral-def$ 
      by auto
    show  $(k, \gamma) \in \{(k'::int, \gamma'). k' * line-integral F$  basis  $\gamma' = 0\}$ 
      using  $zero-line-integral cube-line-integ ass$ 
      by force
  qed
  then have  $ii:(\sum one-chain \in (f' (one-chain1 - \{p. f p = \{\}\})). one-chain-line-integral$ 
 $F$  basis  $one-chain) =$ 
 $(\sum one-chain \in (f' (one-chain1))).$ 
 $one-chain-line-integral F$  basis  $one-chain)$ 
  proof -
    have  $\bigwedge one-chain. one-chain \in (f' (one-chain1)) - (f' (one-chain1 - \{p.$ 
 $f p = \{\}\})) \implies$ 
 $one-chain-line-integral F$  basis
 $one-chain = 0$ 
  proof -
    fix  $one-chain$ 
    assume  $one-chain \in (f' (one-chain1)) - (f' (one-chain1 - \{p. f p =$ 
 $\{\}\}))$ 
    then have  $one-chain = \{\}$ 
      by auto
    then show  $one-chain-line-integral F$  basis  $one-chain = 0$ 
      by (auto simp add:  $one-chain-line-integral-def$ )
  qed
  then have  $0:(\sum one-chain \in f' (one-chain1) - (f' (one-chain1 - \{p. f$ 
 $p = \{\}\})). one-chain-line-integral F$  basis  $one-chain)$ 
 $= 0$ 
    using  $Groups-Big.comm-monoid-add-class.sum.neutral$ 
    by auto
  then have  $(\sum one-chain \in f' (one-chain1). one-chain-line-integral F$  basis
 $one-chain)$ 
 $- (\sum one-chain \in (f' (one-chain1 -$ 
 $\{p. f p = \{\}\})). one-chain-line-integral F$  basis  $one-chain)$ 

```

= 0

proof –
have *finte*: *finite* (*f* ‘ *one-chain1*) **using** *finite-chain1* **by** *auto*
show *?thesis*
using *Groups-Big.sum-diff*[*OF finte*, *of* (*f* ‘ (*one-chain1* – {*p*. *f p* =

{})})
one-chain-line-integral F basis
0
by *auto*
qed
then show ($\sum \text{one-chain} \in (f \text{ ‘ } (one-chain1 - \{p. f p = \{\}\}))$). *one-chain-line-integral F basis one-chain*) =
($\sum \text{one-chain} \in (f \text{ ‘ } (one-chain1))$).
one-chain-line-integral F basis one-chain)
by *auto*
qed
have $\bigwedge (k::int) \gamma. (k, \gamma) \in one-chain1 \implies (f(k, \gamma) = \{\}) \implies f(k, \gamma) \in$
{*chain*. *one-chain-line-integral F basis chain* = 0}
proof–
fix *k::int*
fix *γ::real* $\implies real * real$
assume *ass*: (*k*, *γ*) $\in one-chain1$ (*f* (*k*, *γ*) = {*γ*})
then have *one-chain-line-integral F basis* (*f* (*k*, *γ*)) = 0
using *one-chain-line-integral-def*
by *auto*
then show *f* (*k*, *γ*) $\in \{p'. (one-chain-line-integral F basis p' = 0)\}$
by *auto*
qed
then have *iii*: ($\sum (k::int, \gamma) \in one-chain1 - \{p. f p = \{\}\}$). *k*(line-integral F basis γ)*)
= ($\sum (k::int, \gamma) \in one-chain1$. *k*(line-integral F basis γ)*)
proof–
have $\bigwedge k \gamma. (k, \gamma) \in one-chain1 - (one-chain1 - \{p. f p = \{\}\})$
 $\implies k * (line-integral F basis \gamma) = 0$
proof–
fix *k γ*
assume *ass*: (*k*, *γ*) $\in one-chain1 - (one-chain1 - \{p. f p = \{\}\})$
then have *f*(*k*, *γ*) = {*γ*}
by *auto*
then have *one-chain-line-integral F basis* (*f*(*k*, *γ*)) = 0
by (*auto simp add: one-chain-line-integral-def*)
then have *zero-line-integral:one-chain-line-integral F basis* (*f* (*k*, *γ*)) =
0
using *one-chain-line-integral-def*
by *auto*
then show *k* (line-integral F basis γ)* = 0
using *zero-line-integral cube-line-integ ass*
by *force*

```

qed
then have  $\forall (k::int,\gamma)\in one-chain1 - (one-chain1 - \{p. f p = \{\}\})$ .
   $k*(line-integral F basis \gamma) = 0$  by auto
  then have  $(\sum (k::int,\gamma)\in one-chain1 - (one-chain1 - \{p. f p = \{\}\}))$ .
 $k*(line-integral F basis \gamma) = 0$ 
  using Groups-Big.comm-monoid-add-class.sum.neutral
   $[of one-chain1 - (one-chain1 - \{p. f p = \{\}\}) (\lambda(k::int,\gamma). k*$ 
 $(line-integral F basis \gamma))]$ 
  by (simp add: split-beta)
  then have  $(\sum (k::int,\gamma)\in one-chain1. k*(line-integral F basis \gamma)) -$ 
 $(\sum (k::int,\gamma)\in (one-chain1 - \{p. f p =$ 
 $\{\}\}). k*(line-integral F basis \gamma)) = 0$ 
  using Groups-Big.sum-diff[OF finite-chain1]
  by (metis (no-types) Diff-subset  $\langle(\sum (k, \gamma)\in one-chain1 - (one-chain1$ 
 $- \{p. f p = \{\}\}). k * line-integral F basis \gamma) = 0\rangle \langle\bigwedge f B. B \subseteq one-chain1 \implies$ 
 $sum f (one-chain1 - B) = sum f one-chain1 - sum f B\rangle$ )
  then show ?thesis by auto
qed
show ?thesis using i ii iii by auto
qed
then show ?thesis using one-chain-line-integral-def by auto
qed
show ?thesis using 0 1 by auto
qed
show one-chain-line-integral F basis one-chain1 = one-chain-line-integral F basis
subdiv using 0 1 by auto
qed

```

lemma *chain-reparam-chain'-line-integral-smooth-cubes:*

assumes *chain-reparam-chain' one-chain1 one-chain2*

$\forall (k2,\gamma2)\in one-chain2. \gamma2$ *C1-differentiable-on* $\{0..1\}$

$\forall b\in basis.\forall (k2,\gamma2)\in one-chain2. continuous-on$ (*path-image* $\gamma2$) $(\lambda x. F x \cdot b)$

finite basis

finite one-chain1

boundary-chain one-chain1

boundary-chain one-chain2

$\forall (k,\gamma)\in one-chain1. valid-path \gamma$

shows *one-chain-line-integral F basis one-chain1 = one-chain-line-integral F*
basis one-chain2

$\forall (k, \gamma)\in one-chain1. line-integral-exists F basis \gamma$

proof –

{**fix** *b*

assume $b \in basis$

fix $k \gamma$

assume $(k, \gamma)\in one-chain2$

have *line-integral-exists F* $\{b\} \gamma$

apply (*rule line-integral-exists-smooth*)

using $\langle(k, \gamma) \in one-chain2\rangle$ *assms*(2) **apply** *blast*

using *assms*

using $\langle (k, \gamma) \in \text{one-chain2} \rangle \langle b \in \text{basis} \rangle$ **apply** *blast*
using $\langle b \in \text{basis} \rangle$ **by** *blast*
then have $a: \forall b \in \text{basis}. \forall (k, \gamma) \in \text{one-chain2}. \text{line-integral-exists } F \{b\} \gamma$
by *auto*
have $b: \forall (k2, \gamma2) \in \text{one-chain2}. \text{valid-path } \gamma2$
using *assms(2)*
by (*simp add: C1-differentiable-imp-piecewise case-prod-beta valid-path-def*)

show $\text{one-chain-line-integral } F \text{ basis one-chain1} = \text{one-chain-line-integral } F \text{ basis one-chain2}$
by (*rule chain-reparam-chain'-line-integral[OF assms(1) assms(6) assms(7) a b assms(8) assms(5) assms(4) assms(3)]*)
show $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$
by (*rule chain-reparam-chain'-line-integral[OF assms(1) assms(6) assms(7) a b assms(8) assms(5) assms(4) assms(3)]*)
qed

lemma *chain-subdiv-path-pathimg-subset*:
assumes *chain-subdiv-path* γ *subdiv*
shows $\forall (k', \gamma') \in \text{subdiv}. (\text{path-image } \gamma') \subseteq \text{path-image } \gamma$
using *assms chain-subdiv-path.simps path-image-rec-join*
by (*auto simp add: chain-subdiv-path.simps*)

lemma *reparam-path-image*:
assumes *reparam* $\gamma1$ $\gamma2$
shows $\text{path-image } \gamma1 = \text{path-image } \gamma2$
using *assms*
apply (*auto simp add: reparam-def path-image-def image-def bij-betw-def*)
apply (*force dest!: equalityD2*)
done

lemma *chain-reparam-weak-path-pathimg-subset*:
assumes *chain-reparam-weak-path* γ *subdiv*
shows $\forall (k', \gamma') \in \text{subdiv}. (\text{path-image } \gamma') \subseteq \text{path-image } \gamma$
using *assms*
apply (*auto simp add: chain-reparam-weak-path-def*)
using *path-image-rec-join reparam-path-image* **by** *blast*

lemma *chain-subdiv-chain-pathimg-subset'*:
assumes *chain-subdiv-chain one-chain subdiv*
assumes $(k, \gamma) \in \text{subdiv}$
shows $\exists k' \gamma'. (k', \gamma') \in \text{one-chain} \wedge \text{path-image } \gamma \subseteq \text{path-image } \gamma'$
using *assms unfolding chain-subdiv-chain-def pairwise-def*
apply *auto*
by (*metis chain-subdiv-path.cases coeff-cube-to-path.simps path-image-rec-join path-image-reversepath*)

lemma *chain-subdiv-chain-pathimg-subset*:
assumes *chain-subdiv-chain one-chain subdiv*
shows $\bigcup (\text{path-image } \{ \gamma. \exists k. (k, \gamma) \in \text{subdiv} \}) \subseteq \bigcup (\text{path-image } \{ \gamma. \exists k. (k, \gamma) \in \text{one-chain} \})$

$one-chain\}$)
using *assms* **unfolding** *chain-subdiv-chain-def pairwise-def*
apply *auto*
by (*metis UN-iff assms chain-subdiv-chain-pathimg-subset' subsetCE case-prodI2*)

lemma *chain-reparam-chain'-pathimg-subset'*:
assumes *chain-reparam-chain' one-chain subdiv*
assumes $(k, \gamma) \in subdiv$
shows $\exists k' \gamma'. (k', \gamma') \in one-chain \wedge path-image \gamma \subseteq path-image \gamma'$
using *assms chain-reparam-weak-path-pathimg-subset*
apply (*auto simp add: chain-reparam-chain'-def set-eq-iff*)
using *path-image-reversepath case-prodE case-prodD old.prod.exhaust*
apply (*simp add: list.distinct(1) list.inject rec-join.elims*)
by (*smt case-prodD coeff-cube-to-path.simps rec-join.simps(2) reversepath-simps(2) surj-pair*)

definition *common-reparam-exists:: (int \times (real \Rightarrow real \times real)) set \Rightarrow (int \times (real \Rightarrow real \times real)) set \Rightarrow bool* **where**
 $common-reparam-exists\ one-chain1\ one-chain2 \equiv$
 $(\exists subdiv\ ps1\ ps2.$
 $chain-reparam-chain'\ (one-chain1 - ps1)\ subdiv \wedge$
 $chain-reparam-chain'\ (one-chain2 - ps2)\ subdiv \wedge$
 $(\forall (k, \gamma) \in subdiv. \gamma\ C1-differentiable-on\ \{0..1\}) \wedge$
 $boundary-chain\ subdiv \wedge$
 $(\forall (k, \gamma) \in ps1. point-path\ \gamma) \wedge$
 $(\forall (k, \gamma) \in ps2. point-path\ \gamma))$

lemma *common-reparam-exists-imp-eq-line-integral:*
assumes *finite-basis: finite basis and*
 $finite\ one-chain1$
 $finite\ one-chain2$
 $boundary-chain\ (one-chain1::(int \times (real \Rightarrow real \times real))\ set)$
 $boundary-chain\ (one-chain2::(int \times (real \Rightarrow real \times real))\ set)$
 $\forall (k2, \gamma2) \in one-chain2. \forall b \in basis. continuous-on\ (path-image\ \gamma2)\ (\lambda x. F\ x \cdot$
b)
 $(common-reparam-exists\ one-chain1\ one-chain2)$
 $(\forall (k, \gamma) \in one-chain1. valid-path\ \gamma)$
 $(\forall (k, \gamma) \in one-chain2. valid-path\ \gamma)$
shows $one-chain-line-integral\ F\ basis\ one-chain1 = one-chain-line-integral\ F$
 $basis\ one-chain2$
 $\forall (k, \gamma) \in one-chain1. line-integral-exists\ F\ basis\ \gamma$

proof –

obtain *subdiv ps1 ps2 where props:*
 $chain-reparam-chain'\ (one-chain1 - ps1)\ subdiv$
 $chain-reparam-chain'\ (one-chain2 - ps2)\ subdiv$
 $(\forall (k, \gamma) \in subdiv. \gamma\ C1-differentiable-on\ \{0..1\})$
 $boundary-chain\ subdiv$
 $(\forall (k, \gamma) \in ps1. point-path\ \gamma)$


```

    (∀(k, γ)∈ps2. point-path γ)
  using assms
  by (auto simp add: common-reparam-exists-def)
  have subdiv-valid: (∀(k, γ)∈subdiv. valid-path γ)
  apply (simp add: valid-path-def)
  using props(3)
  using C1-differentiable-imp-piecewise by blast
  have onechain-boundary1: boundary-chain (one-chain1 - ps1) using assms(4)
  by (auto simp add: boundary-chain-def)
  have onechain-boundary2: boundary-chain (one-chain2 - ps1) using assms(5)
  by (auto simp add: boundary-chain-def)
  {fix k2 γ2 b
    assume ass: (k2, γ2)∈subdiv b∈basis
    have ∧ k γ. (k, γ) ∈ subdiv ⇒ ∃ k' γ'. (k', γ') ∈ one-chain2 ∧ path-image γ
    ⊆ path-image γ'
      by (meson chain-reparam-chain'-pathimg-subset' props Diff-subset subsetCE)
    then have continuous-on (path-image γ2) (λx. F x · b)
      using assms(6) continuous-on-subset[where ?f = (λx. F x · b)] ass
      apply (auto simp add: subset-iff)
      by (metis (mono-tags, lifting) case-prodD)}
  then have cont-field: ∀ b∈basis. ∀(k2, γ2)∈subdiv. continuous-on (path-image
  γ2) (λx. F x · b)
  by auto
  have one-chain1-ps-valid: (∀(k, γ)∈one-chain1 - ps1. valid-path γ) using assms
  by auto
  have one-chain2-ps-valid: (∀(k, γ)∈one-chain2 - ps1. valid-path γ) using assms
  by auto
  have 0: one-chain-line-integral F basis (one-chain1 - ps1) = one-chain-line-integral
  F basis subdiv
  apply (rule chain-reparam-chain'-line-integral-smooth-cubes[OF props(1) props(3)
  cont-field finite-basis])
  using props assms
  apply blast
  using props assms
  using onechain-boundary1 apply blast
  using props assms
  apply blast
  using one-chain1-ps-valid by blast
  have 1: ∀(k, γ)∈(one-chain1 - ps1). line-integral-exists F basis γ
  apply (rule chain-reparam-chain'-line-integral-smooth-cubes[OF props(1) props(3)
  cont-field finite-basis])
  using props assms
  apply blast
  using props assms
  using onechain-boundary1 apply blast
  using props assms
  apply blast
  using one-chain1-ps-valid by blast
  have 2: one-chain-line-integral F basis (one-chain2 - ps2) = one-chain-line-integral

```

```

F basis subdiv
  apply(rule chain-reparam-chain'-line-integral-smooth-cubes[OF props(2) props(3)
cont-field finite-basis])
  using props assms
  apply blast
  apply (simp add: assms(5) boundary-chain-diff)
  apply (simp add: props(4))
  by (simp add: assms(9))
  have 3:  $\forall (k, \gamma) \in (\text{one-chain2} - \text{ps2}). \text{line-integral-exists } F \text{ basis } \gamma$ 
  apply(rule chain-reparam-chain'-line-integral-smooth-cubes[OF props(2) props(3)
cont-field finite-basis])
  using props assms
  apply blast
  apply (simp add: assms(5) boundary-chain-diff)
  apply (simp add: props(4))
  by (simp add: assms(9))
  show line-int-ex-chain1:  $\forall (k, \gamma) \in \text{one-chain1}. \text{line-integral-exists } F \text{ basis } \gamma$ 
  using 0
  using 1 finite-basis line-integral-exists-point-path props(5) by fastforce
  then show one-chain-line-integral F basis one-chain1 = one-chain-line-integral
F basis one-chain2
  using 0 1 2 3
  using assms(2-3) finite-basis one-chain-line-integral-point-paths props(5) props(6)
by auto
qed

```

```

definition subcube :: real  $\Rightarrow$  real  $\Rightarrow$  (int  $\times$  (real  $\Rightarrow$  real  $\times$  real))  $\Rightarrow$  (int  $\times$  (real
 $\Rightarrow$  real  $\times$  real)) where
  subcube a b cube = (fst cube, subpath a b (snd cube))

```

```

lemma subcube-valid-path:
  assumes valid-path (snd cube) a  $\in$  {0..1} b  $\in$  {0..1}
  shows valid-path (snd (subcube a b cube))
  using valid-path-subpath[OF assms] by (auto simp add: subcube-def)

```

```

end
theory Green
  imports Paths Derivs Integrals General-Utills
begin

```

```

lemma frontier-Un-subset-Un-frontier:
  frontier (s  $\cup$  t)  $\subseteq$  (frontier s)  $\cup$  (frontier t)
  by (simp add: frontier-def Un-Diff) (auto simp add: closure-def interior-def is-
limpt-def)

```

```

definition has-partial-derivative:: ('a::euclidean-space)  $\Rightarrow$  'b::euclidean-space)  $\Rightarrow$ 
'a  $\Rightarrow$  ('a  $\Rightarrow$  'b)  $\Rightarrow$  ('a)  $\Rightarrow$  bool where
  has-partial-derivative F base-vec F' a
     $\equiv$  (( $\lambda x::'a::euclidean-space. F( a - ((a \cdot \text{base-vec}) *_R \text{base-vec})) + (x \cdot$ 

```

$\text{base-vec}) *_{\mathbb{R}} \text{base-vec}))$
 $\text{has-derivative } F') \text{ (at } a)$

definition *has-partial-vector-derivative*:: $(('a::\text{euclidean-space}) \Rightarrow 'b::\text{euclidean-space})$
 $\Rightarrow 'a \Rightarrow ('b) \Rightarrow ('a) \Rightarrow \text{bool}$ **where**
 $\text{has-partial-vector-derivative } F \text{ base-vec } F' a$
 $\equiv ((\lambda x. F(a - ((a \cdot \text{base-vec}) *_{\mathbb{R}} \text{base-vec})) + x *_{\mathbb{R}} \text{base-vec}))$
 $\text{has-vector-derivative } F') \text{ (at } (a \cdot \text{base-vec}))$

definition *partially-vector-differentiable* **where**
 $\text{partially-vector-differentiable } F \text{ base-vec } p \equiv (\exists F'. \text{has-partial-vector-derivative } F$
 $\text{base-vec } F' p)$

definition *partial-vector-derivative*:: $(('a::\text{euclidean-space}) \Rightarrow 'b::\text{euclidean-space})$
 $\Rightarrow 'a \Rightarrow 'a \Rightarrow 'b$ **where**
 $\text{partial-vector-derivative } F \text{ base-vec } a$
 $\equiv (\text{vector-derivative } (\lambda x. F(a - ((a \cdot \text{base-vec}) *_{\mathbb{R}} \text{base-vec})) + x *_{\mathbb{R}}$
 $\text{base-vec})) \text{ (at } (a \cdot \text{base-vec})))$

lemma *partial-vector-derivative-works*:

assumes $\text{partially-vector-differentiable } F \text{ base-vec } a$

shows $\text{has-partial-vector-derivative } F \text{ base-vec } (\text{partial-vector-derivative } F \text{ base-vec}$
 $a) a$

proof –

obtain F' **where** $F'\text{-prop}$: $((\lambda x. F(a - ((a \cdot \text{base-vec}) *_{\mathbb{R}} \text{base-vec})) + x *_{\mathbb{R}}$
 $\text{base-vec}))$

$\text{has-vector-derivative } F') \text{ (at } (a \cdot \text{base-vec}))$

using assms $\text{partially-vector-differentiable-def}$ $\text{has-partial-vector-derivative-def}$

by blast

show $?thesis$

using $\text{Derivative.differentiableI-vector}[OF F'\text{-prop}]$

by $(\text{simp add: vector-derivative-works partial-vector-derivative-def[symmetric]}$
 $\text{has-partial-vector-derivative-def[symmetric]})$

qed

lemma *fundamental-theorem-of-calculus-partial-vector*:

fixes $a b::\text{real}$ **and**

$F::('a::\text{euclidean-space} \Rightarrow 'b::\text{euclidean-space})$ **and**

$i:: 'a$ **and**

$j:: 'b$ **and**

$F\text{-j-i}::('a::\text{euclidean-space} \Rightarrow \text{real})$

assumes $a\text{-leq-}b$: $a \leq b$ **and**

Base-vecs : $i \in \text{Basis } j \in \text{Basis}$ **and**

no-i-component : $c \cdot i = 0$ **and**

has-partial-deriv : $\forall p \in D. \text{has-partial-vector-derivative } (\lambda x. (F x) \cdot j) i (F\text{-j-i}$
 $p) p$ **and**

$\text{domain-subset-of-}D$: $\{x *_{\mathbb{R}} i + c \mid x. a \leq x \wedge x \leq b\} \subseteq D$

shows $((\lambda x. F\text{-j-i}(x *_{\mathbb{R}} i + c)) \text{has-integral}$

$F(b *_{\mathbb{R}} i + c) \cdot j - F(a *_{\mathbb{R}} i + c) \cdot j) (cbox a b)$

proof –

let $?domain = \{v. \exists x. a \leq x \wedge x \leq b \wedge v = x *_R i + c\}$
have $\forall x \in ?domain. \text{has-partial-vector-derivative } (\lambda x. (F x) \cdot j) i (F\text{-}j\text{-}i x) x$
using *has-partial-deriv domain-subset-of-D*
by *auto*
then have $\forall x \in (cbox\ a\ b). ((\lambda x. F(x *_R i + c) \cdot j) \text{has-vector-derivative } (F\text{-}j\text{-}i(x *_R i + c))) (at\ x)$
proof(*clarsimp simp add: has-partial-vector-derivative-def*)
fix $x::real$
assume *range-of-x: $a \leq x \leq b$*
assume *ass2: $\forall x. (\exists z \geq a. z \leq b \wedge x = z *_R i + c) \longrightarrow$*
 $((\lambda z. F(x - (x \cdot i) *_R i + z *_R i) \cdot j) \text{has-vector-derivative } F\text{-}j\text{-}i$
 $x) (at\ (x \cdot i))$
have $((\lambda z. F((x *_R i + c) - ((x *_R i + c) \cdot i) *_R i + z *_R i) \cdot j)$
 $\text{has-vector-derivative } F\text{-}j\text{-}i(x *_R i + c)) (at\ ((x *_R i + c) \cdot i))$
using *range-of-x ass2 by auto*
then have $0: ((\lambda x. F(c + x *_R i) \cdot j) \text{has-vector-derivative } F\text{-}j\text{-}i(x *_R i +$
 $c)) (at\ x)$
by (*simp add: assms(2) inner-left-distrib no-i-component*)
have $1: (\lambda x. F(x *_R i + c) \cdot j) = (\lambda x. F(c + x *_R i) \cdot j)$
by (*simp add: add commute*)
then show $((\lambda x. F(x *_R i + c) \cdot j) \text{has-vector-derivative } F\text{-}j\text{-}i(x *_R i + c))$
 $(at\ x)$
using 0 and 1 **by** *auto*
qed
then have $\forall x \in (cbox\ a\ b). ((\lambda x. F(x *_R i + c) \cdot j) \text{has-vector-derivative } (F\text{-}j\text{-}i(x *_R i + c))) (at\ \text{within } x\ (cbox\ a\ b))$
using *has-vector-derivative-at-within*
by *blast*
then show $(\lambda x. F\text{-}j\text{-}i(x *_R i + c)) \text{has-integral}$
 $F(b *_R i + c) \cdot j - F(a *_R i + c) \cdot j (cbox\ a\ b)$
using *fundamental-theorem-of-calculus[of a b ($\lambda x::real. F(x *_R i + c) \cdot j$)*
 $(\lambda x::real. F\text{-}j\text{-}i(x *_R i + c))]$ **and**
 $a\text{-leq-}b$
by *auto*
qed

lemma *fundamental-theorem-of-calculus-partial-vector-gen:*

fixes $k1\ k2::real$ **and**

$F::('a::euclidean-space \Rightarrow 'b::euclidean-space)$ **and**

$i::'a$ **and**

$F\text{-}i::('a::euclidean-space \Rightarrow 'b)$

assumes $a\text{-leq-}b: k1 \leq k2$ **and**

$unit\text{-}len: i \cdot i = 1$ **and**

$no\text{-}i\text{-}component: c \cdot i = 0$ **and**

$has\text{-}partial\text{-}deriv: \forall p \in D. \text{has-partial-vector-derivative } F\ i\ (F\text{-}i\ p)\ p$ **and**

$domain\text{-}subset\text{-}of\text{-}D: \{v. \exists x. k1 \leq x \wedge x \leq k2 \wedge v = x *_R i + c\} \subseteq D$

shows $(\lambda x. F\text{-}i(x *_R i + c)) \text{has-integral}$

$F(k2 *_R i + c) - F(k1 *_R i + c) (cbox\ k1\ k2)$

proof –

let $?domain = \{v. \exists x. k1 \leq x \wedge x \leq k2 \wedge v = x *_R i + c\}$

have $\forall x \in ?domain. \text{has-partial-vector-derivative } F \ i \ (F-i \ x) \ x$

using *has-partial-deriv domain-subset-of-D*

by *auto*

then have $\forall x \in (cbox \ k1 \ k2). ((\lambda x. F(x *_R i + c)) \text{has-vector-derivative } (F-i(x *_R i + c))) \ (at \ x)$

proof (*clarsimp simp add: has-partial-vector-derivative-def*)

fix $x::real$

assume *range-of-x: $k1 \leq x \leq k2$*

assume *ass2: $\forall x. (\exists z \geq k1. z \leq k2 \wedge x = z *_R i + c) \longrightarrow ((\lambda z. F(x - (x \cdot i) *_R i + z *_R i)) \text{has-vector-derivative } F-i \ x)$*

(*at* $(x \cdot i)$)

have $((\lambda z. F((x *_R i + c) - ((x *_R i + c) \cdot i) *_R i + z *_R i)) \text{has-vector-derivative } F-i \ (x *_R i + c)) \ (at \ ((x *_R i + c) \cdot i))$

using *range-of-x ass2 by auto*

then have $0: ((\lambda x. F(c + x *_R i)) \text{has-vector-derivative } F-i \ (x *_R i + c)) \ (at \ x)$

by (*simp add: inner-commute inner-right-distrib no-i-component unit-len*)

have $1: (\lambda x. F(x *_R i + c)) = (\lambda x. F(c + x *_R i))$

by (*simp add: add.commute*)

then show $((\lambda x. F(x *_R i + c)) \text{has-vector-derivative } F-i \ (x *_R i + c)) \ (at \ x)$

using 0 and 1 **by** *auto*

qed

then have $\forall x \in (cbox \ k1 \ k2). ((\lambda x. F(x *_R i + c)) \text{has-vector-derivative } (F-i(x *_R i + c))) \ (at-within \ x \ (cbox \ k1 \ k2))$

using *has-vector-derivative-at-within*

by *blast*

then show $(\lambda x. F-i(x *_R i + c)) \text{has-integral}$

$$F(k2 *_R i + c) - F(k1 *_R i + c) \ (cbox \ k1 \ k2)$$

using *fundamental-theorem-of-calculus[of k1 k2 ($\lambda x::real. F(x *_R i + c)$)]* **and**

a-leq-b

by *auto*

qed

lemma *add-scale-img:*

assumes $a < b$ **shows** $(\lambda x::real. a + (b - a) * x) \ ' \{0 .. 1\} = \{a .. b\}$

using *assms*

apply (*auto simp: algebra-simps affine-ineq image-iff*)

using *less-eq-real-def* **apply** *force*

apply (*rule-tac x=(x-a)/(b-a) in bexI*)

apply (*auto simp: field-simps*)

done

lemma *add-scale-img':*

assumes $a \leq b$

shows $(\lambda x::real. a + (b - a) * x) \ ' \{0 .. 1\} = \{a .. b\}$

proof (*cases a = b*)

```

case True
then show ?thesis by force
next
case False
then show ?thesis
using add-scale-imp assms by auto
qed

```

definition *analytically-valid*:: *'a::euclidean-space set* \Rightarrow (*'a* \Rightarrow *'b::\{euclidean-space,times,zero-neq-one\}*) \Rightarrow *'a* \Rightarrow *bool* **where**
analytically-valid s F i \equiv
 $(\forall a \in s. \text{partially-vector-differentiable } F \ i \ a) \wedge$
 $\text{continuous-on } s \ F \wedge$ — **TODO**: should we replace this with saying that F is
partially differentiable on Dy ,
— i.e. there is a partial derivative on every dimension
 $\text{integrable lborel } (\lambda p. (\text{partial-vector-derivative } F \ i) \ p * \text{indicator } s \ p) \wedge$
 $(\lambda x. \text{integral UNIV } (\lambda y. (\text{partial-vector-derivative } F \ i) (y *_R i + x *_R (\sum b$
 $\in (\text{Basis} - \{i\}). b)))$
 $* (\text{indicator } s (y *_R i + x *_R (\sum b \in \text{Basis} - \{i\}. b)))) \in \text{borel-measurable}$
lborel

lemma *analytically-valid-imp-part-deriv-integrable-on*:
assumes *analytically-valid (s::(real*real) set) (f::(real*real) \Rightarrow real) i*
shows *(partial-vector-derivative f i) integrable-on s*
proof —
have *integrable lborel* $(\lambda p. \text{partial-vector-derivative } f \ i \ p * \text{indicator } s \ p)$
using *assms*
by *(simp add: analytically-valid-def indic-ident)*
then have *integrable lborel* $(\lambda p::(\text{real*real}). \text{if } p \in s \text{ then } \text{partial-vector-derivative}$
 $f \ i \ p \text{ else } 0)$
using *indic-ident[of partial-vector-derivative f i]*
by *(simp add: indic-ident)*
then have $(\lambda x. \text{if } x \in s \text{ then } \text{partial-vector-derivative } f \ i \ x \text{ else } 0)$ *integrable-on*
UNIV
using *Equivalence-Lebesgue-Henstock-Integration.integrable-on-lborel*
by auto
then show *(partial-vector-derivative f i) integrable-on s*
using *integrable-restrict-UNIV*
by auto
qed

definition *typeII-twoCube* :: $((\text{real} * \text{real}) \Rightarrow (\text{real} * \text{real})) \Rightarrow \text{bool}$ **where**
typeII-twoCube twoC
 $\equiv \exists a \ b \ g1 \ g2. a < b \wedge (\forall x \in \{a..b\}. g2 \ x \leq g1 \ x) \wedge$
 $\text{twoC} = (\lambda(y, x). ((1 - y) * (g2 ((1-x)*a + x*b)) + y * (g1$

$((1-x)*a + x*b)),$
 $(1-x)*a + x*b)) \wedge$
 $g1 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\} \wedge$
 $g2 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\}$

abbreviation *unit-cube* **where** *unit-cube* \equiv *cbox* (0,0) (1::real,1::real)

definition *cubeImage*:: *two-cube* \Rightarrow ((*real***real*) *set*) **where**
cubeImage twoC \equiv (*twoC* ' *unit-cube*)

lemma *typeII-twoCubeImg*:

assumes *typeII-twoCube twoC*

shows $\exists a b g1 g2. a < b \wedge (\forall x \in \{a .. b\}. g2 x \leq g1 x) \wedge$

$\text{cubeImage twoC} = \{(y,x). x \in \{a..b\} \wedge y \in \{g2 x .. g1 x\}\}$

$\wedge \text{twoC} = (\lambda(y, x). ((1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b), (1 - x) * a + x * b))$

$\wedge g1 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\} \wedge g2 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\}$

using *assms*

proof (*simp add: typeII-twoCube-def cubeImage-def image-def*)

assume $\exists a b. a < b \wedge (\exists g1 g2. (\forall x \in \{a..b\}. g2 x \leq g1 x) \wedge$

$\text{twoC} = (\lambda(y, x). ((1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b), (1 - x) * a + x * b)) \wedge$

$g1 \text{ piecewise-}C1\text{-differentiable-on } \{a..b\} \wedge g2 \text{ piecewise-}C1\text{-differentiable-on } \{a..b\})$

then obtain *a b g1 g2* **where**

twoCisTypeII: $a < b$

$(\forall x \in \{a..b\}. g2 x \leq g1 x)$

$\text{twoC} = (\lambda(y, x). ((1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b), (1 - x) * a + x * b))$

$g1 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\}$

$g2 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\}$

by *auto*

have *ab1*: $a - x * a + x * b \leq b$ **if** *a1*: $0 \leq x \leq 1$ **for** *x*

using *that* **apply** (*simp add: algebra-simps*)

by (*metis affine-ineq less-eq-real-def mult.commute twoCisTypeII(1)*)

have *ex*: $\exists z \in \text{Green.unit-cube}.$

$(d, c) =$

(*case z of*

$(y, x) \Rightarrow$

$(g2 (a - x * a + x * b) - y * g2 (a - x * a + x * b) + y * g1 (a - x * a + x * b),$

$a - x * a + x * b))$

if *c-bounds*: $a \leq c \leq b$ **and** *d-bounds*: $g2 c \leq d \leq g1 c$ **for** *c d*

proof -

have *b-minus-a-nzero*: $b - a \neq 0$ **using** *twoCisTypeII(1)* **by** *auto*

have *x-witness*: $\exists k1. c = (1 - k1)*a + k1 * b \wedge 0 \leq k1 \wedge k1 \leq 1$

apply (*rule-tac x=(c - a)/(b - a) in exI*)

using *c-bounds* $\langle a < b \rangle$ **apply** (*simp add: divide-simps algebra-simps*)

```

    using sum-sqs-eq by blast
  have y-witness:  $\exists k2. d = (1 - k2) * (g2 c) + k2 * (g1 c) \wedge 0 \leq k2 \wedge k2 \leq 1$ 
  proof (cases  $g1 c - g2 c = 0$ )
    case True
      with d-bounds show ?thesis by (fastforce simp add: algebra-simps)
    next
      case False
        let ?k2 =  $(d - g2 c) / (g1 c - g2 c)$ 
        have k2-in-bounds:  $0 \leq ?k2 \wedge ?k2 \leq 1$ 
          using twoCisTypeII(2) c-bounds d-bounds False by simp
        have d =  $(1 - ?k2) * g2 c + ?k2 * g1 c$ 
          using False sum-sqs-eq by (fastforce simp add: divide-simps algebra-simps)
        with k2-in-bounds show ?thesis
          by fastforce
      qed
    show  $\exists x \in \text{unit-cube}. (d, c) = (\text{case } x \text{ of } (y, x) \Rightarrow (g2 (a - x * a + x * b) - y * g2 (a - x * a + x * b) + y * g1 (a - x * a + x * b), a - x * a + x * b))$ 
      using x-witness y-witness by (force simp add: left-diff-distrib)
    qed
  have  $\{y. \exists x \in \text{unit-cube}. y = \text{twoC } x\} = \{(y, x). a \leq x \wedge x \leq b \wedge g2 x \leq y \wedge y \leq g1 x\}$ 
    apply (auto simp add: twoCisTypeII ab1 left-diff-distrib ex)
    using order.order-iff-strict twoCisTypeII(1) apply fastforce
    apply (smt affine-ineq atLeastAtMost-iff mult-left-mono twoCisTypeII)+
    done
  then show  $\exists a b. a < b \wedge (\exists g1 g2. (\forall x \in \{a..b\}. g2 x \leq g1 x) \wedge \{y. \exists x \in \text{unit-cube}. y = \text{twoC } x\} = \{(y, x). a \leq x \wedge x \leq b \wedge g2 x \leq y \wedge y \leq g1 x\} \wedge \text{twoC} = (\lambda(y, x). ((1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b), (1 - x) * a + x * b)) \wedge g1 \text{ piecewise-}C1\text{-differentiable-on } \{a..b\} \wedge g2 \text{ piecewise-}C1\text{-differentiable-on } \{a..b\})$ 
    using twoCisTypeII by blast
  qed

```

definition *horizontal-boundary* :: *two-cube* \Rightarrow *one-chain* **where**
horizontal-boundary twoC $\equiv \{(1, (\lambda x. \text{twoC}(x, 0))), (-1, (\lambda x. \text{twoC}(x, 1)))\}$

definition *vertical-boundary* :: *two-cube* \Rightarrow *one-chain* **where**
vertical-boundary twoC $\equiv \{(-1, (\lambda y. \text{twoC}(0, y))), (1, (\lambda y. \text{twoC}(1, y)))\}$

definition *boundary* :: *two-cube* \Rightarrow *one-chain* **where**
boundary twoC $\equiv \text{horizontal-boundary twoC} \cup \text{vertical-boundary twoC}$

definition *valid-two-cube* **where**
valid-two-cube twoC $\equiv \text{card } (\text{boundary twoC}) = 4$

definition *two-chain-integral*:: *two-chain* \Rightarrow $((\text{real} * \text{real}) \Rightarrow (\text{real})) \Rightarrow \text{real}$ **where**
two-chain-integral twoChain F $\equiv \sum_{C \in \text{twoChain}. (\text{integral } (\text{cubeImage } C) F)$

definition *valid-two-chain where*

$valid-two-chain\ twoChain \equiv (\forall\ twoCube \in\ twoChain.\ valid-two-cube\ twoCube) \wedge pairwise\ (\lambda c1\ c2.\ ((boundary\ c1) \cap (boundary\ c2)) = \{\})\ twoChain \wedge inj-on\ cubeImage\ twoChain$

definition *two-chain-boundary:: two-chain \Rightarrow one-chain where*

$two-chain-boundary\ twoChain == \bigcup (boundary\ 'twoChain)$

definition *gen-division where*

$gen-division\ s \equiv (finite\ S \wedge (\bigcup S = s) \wedge pairwise\ (\lambda X\ Y.\ negligible\ (X \cap Y))\ S)$

definition *two-chain-horizontal-boundary:: two-chain \Rightarrow one-chain where*

$two-chain-horizontal-boundary\ twoChain \equiv \bigcup (horizontal-boundary\ 'twoChain)$

definition *two-chain-vertical-boundary:: two-chain \Rightarrow one-chain where*

$two-chain-vertical-boundary\ twoChain \equiv \bigcup (vertical-boundary\ 'twoChain)$

definition *only-horizontal-division where*

only-horizontal-division one-chain two-chain

$\equiv \exists\ \mathcal{H}\ \mathcal{V}.\ finite\ \mathcal{H} \wedge finite\ \mathcal{V} \wedge$

$(\forall (k, \gamma) \in \mathcal{H}.$

$(\exists (k', \gamma') \in two-chain-horizontal-boundary\ two-chain.$

$(\exists a \in \{0..1\}.\ \exists b \in \{0..1\}.\ a \leq b \wedge subpath\ a\ b\ \gamma' = \gamma))) \wedge$

$(common-sudiv-exists\ (two-chain-vertical-boundary\ two-chain)\ \mathcal{V}$

$\vee\ common-reparam-exists\ \mathcal{V}\ (two-chain-vertical-boundary\ two-chain))$

\wedge

$boundary-chain\ \mathcal{V} \wedge$

$one-chain = \mathcal{H} \cup \mathcal{V} \wedge (\forall (k, \gamma) \in \mathcal{V}.\ valid-path\ \gamma)$

lemma *sum-zero-set:*

assumes $\forall x \in s.\ f\ x = 0\ finite\ s\ finite\ t$

shows $sum\ f\ (s \cup t) = sum\ f\ t$

using *assms*

by (*simp add: IntE sum.union-inter-neutral sup-commute*)

abbreviation *valid-typeII-division s twoChain $\equiv ((\forall\ twoCube \in\ twoChain.\ typeII-twoCube\ twoCube) \wedge$*

(gen-division s (cubeImage 'twoChain)) \wedge

(valid-two-chain twoChain))

lemma *two-chain-vertical-boundary-is-boundary-chain:*

shows *boundary-chain (two-chain-vertical-boundary twoChain)*

by (*simp add: boundary-chain-def two-chain-vertical-boundary-def vertical-boundary-def*)

lemma *two-chain-horizontal-boundary-is-boundary-chain:*

shows *boundary-chain (two-chain-horizontal-boundary twoChain)*

by(*simp add: boundary-chain-def two-chain-horizontal-boundary-def horizontal-boundary-def*)

definition *typeI-twoCube* :: *two-cube* \Rightarrow *bool* **where**

typeI-twoCube (*twoC*::*two-cube*)
 $\equiv \exists a b g1 g2. a < b \wedge (\forall x \in \{a..b\}. g2 x \leq g1 x) \wedge$
 $twoC = (\lambda(x,y). ((1-x)*a + x*b,$
 $(1 - y) * (g2 ((1-x)*a + x*b)) + y * (g1$
 $((1-x)*a + x*b)))) \wedge$
 $g1 \text{ piecewise-C1-differentiable-on } \{a..b\} \wedge$
 $g2 \text{ piecewise-C1-differentiable-on } \{a..b\}$

lemma *typeI-twoCubeImg*:

assumes *typeI-twoCube twoC*

shows $\exists a b g1 g2. a < b \wedge (\forall x \in \{a .. b\}. g2 x \leq g1 x) \wedge$
 $cubeImage \ twoC = \{(x,y). x \in \{a..b\} \wedge y \in \{g2 x .. g1 x\}\} \wedge$
 $twoC = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2 ((1 - x) * a$
 $+ x * b) + y * g1 ((1 - x) * a + x * b))) \wedge$
 $g1 \text{ piecewise-C1-differentiable-on } \{a .. b\} \wedge g2 \text{ piecewise-C1-differentiable-on } \{a .. b\}$

proof –

have $\exists a b. a < b \wedge$
 $(\exists g1 g2. (\forall x \in \{a..b\}. g2 x \leq g1 x) \wedge$
 $twoC = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2 ((1 - x) * a$
 $+ x * b) + y * g1 ((1 - x) * a + x * b))) \wedge$
 $g1 \text{ piecewise-C1-differentiable-on } \{a .. b\} \wedge g2 \text{ piecewise-C1-differentiable-on } \{a .. b\})$

using *assms* **by** (*simp add: typeI-twoCube-def*)

then obtain *a b g1 g2* **where**

twoCisTypeI: $a < b$
 $(\forall x \in \{a..b\}. g2 x \leq g1 x)$
 $twoC = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b)))$
 $g1 \text{ piecewise-C1-differentiable-on } \{a .. b\}$
 $g2 \text{ piecewise-C1-differentiable-on } \{a .. b\}$

by *auto*

have *ex*: $\exists z \in Green.unit-cube.$

$(c, d) =$
 $(case \ z \ of$
 $(x, y) \Rightarrow$
 $(a - x * a + x * b,$
 $g2 (a - x * a + x * b) - y * g2 (a - x * a + x * b) + y * g1 (a$
 $- x * a + x * b)))$

if *c-bounds*: $a \leq c \leq b$ **and** *d-bounds*: $g2 c \leq d \leq g1 c$ **for** *c d*

proof –

have *x-witness*: $\exists k1. c = (1 - k1)*a + k1 * b \wedge 0 \leq k1 \wedge k1 \leq 1$

proof –

let *?k1* = $(c - a)/(b - a)$

have *k1-in-bounds*: $0 \leq (c - a)/(b - a) \wedge (c - a)/(b - a) \leq 1$

using *twoCisTypeI(1)* *c-bounds* **by** *simp*

```

have  $c = (1 - ?k1)*a + ?k1 * b$ 
  using twoCisTypeI(1) sum-sqs-eq
  by (auto simp add: divide-simps algebra-simps)
then show ?thesis
  using twoCisTypeI k1-in-bounds by fastforce
qed
have y-witness:  $\exists k2. d = (1 - k2)*(g2 c) + k2 * (g1 c) \wedge 0 \leq k2 \wedge k2 \leq 1$ 
proof (cases  $g1 c - g2 c = 0$ )
  case True
    with d-bounds show ?thesis
    by force
  next
    case False
    let  $?k2 = (d - g2 c)/(g1 c - g2 c)$ 
    have k2-in-bounds:  $0 \leq ?k2 \wedge ?k2 \leq 1$  using twoCisTypeI(2) c-bounds
    d-bounds False by simp
    have  $d = (1 - ?k2) * g2 c + ?k2 * g1 c$ 
    using False apply (simp add: divide-simps algebra-simps)
    using sum-sqs-eq by fastforce
    then show ?thesis using k2-in-bounds by fastforce
  qed
show  $\exists x \in \text{unit-cube}. (c, d) =$ 
  (case  $x$  of  $(x, y) \Rightarrow (a - x * a + x * b, g2 (a - x * a + x * b) - y * g2 (a - x * a + x * b) + y * g1 (a - x * a + x * b))$ )
  using x-witness y-witness by (force simp add: left-diff-distrib)
qed
have  $\{y. \exists x \in \text{unit-cube}. y = \text{twoC } x\} = \{(x, y). a \leq x \wedge x \leq b \wedge g2 x \leq y \wedge y \leq g1 x\}$ 
apply (auto simp add: twoCisTypeI left-diff-distrib ex)
using less-eq-real-def twoCisTypeI(1) apply auto[1]
apply (smt affine-ineq twoCisTypeI)
apply (smt affine-ineq atLeastAtMost-iff mult-left-mono twoCisTypeI)+
done
then show ?thesis
unfolding cubeImage-def image-def using twoCisTypeI by auto
qed

```

lemma *typeI-cube-explicit-spec*:

assumes *typeI-twoCube twoC*

```

shows  $\exists a b g1 g2. a < b \wedge (\forall x \in \{a .. b\}. g2 x \leq g1 x) \wedge$ 
   $\text{cubeImage } \text{twoC} = \{(x, y). x \in \{a..b\} \wedge y \in \{g2 x .. g1 x\}\}$ 
   $\wedge \text{twoC} = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b)))$ 
   $\wedge g1 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\} \wedge g2 \text{ piecewise-}C1\text{-differentiable-on } \{a .. b\}$ 
   $\wedge (\lambda x. \text{twoC}(x, 0)) = (\lambda x. (a + (b - a) * x, g2 (a + (b - a) * x)))$ 
   $\wedge (\lambda y. \text{twoC}(1, y)) = (\lambda x. (b, g2 b + x *_R (g1 b - g2 b)))$ 
   $\wedge (\lambda x. \text{twoC}(x, 1)) = (\lambda x. (a + (b - a) * x, g1 (a + (b - a) * x)))$ 

```

$x)))$

$$\wedge (\lambda y. \text{twoC}(0, y)) = (\lambda x. (a, g2\ a + x *_R (g1\ a - g2\ a)))$$

proof –

let $?bottom\text{-}edge = (\lambda x. \text{twoC}(x, 0))$
let $?right\text{-}edge = (\lambda y. \text{twoC}(1, y))$
let $?top\text{-}edge = (\lambda x. \text{twoC}(x, 1))$
let $?left\text{-}edge = (\lambda y. \text{twoC}(0, y))$
obtain $a\ b\ g1\ g2$ **where**
 $\text{twoCisTypeI}: a < b$
 $(\forall x \in \text{cbox } a\ b. g2\ x \leq g1\ x)$
 $\text{cubeImage } \text{twoC} = \{(x,y). x \in \text{cbox } a\ b \wedge y \in \text{cbox } (g2\ x)\ (g1\ x)\}$
 $\text{twoC} = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) * a + x * b)))$
 $g1\ \text{piecewise-C1-differentiable-on } \{a .. b\}$
 $g2\ \text{piecewise-C1-differentiable-on } \{a .. b\}$
using assms **and** $\text{typeI-twoCubeImg[oftwoC]}$ **by** auto
have $\text{bottom-edge-explicit}: ?bottom\text{-}edge = (\lambda x. (a + (b - a) * x, g2\ (a + (b - a) * x)))$
by $(\text{simp add: twoCisTypeI}(4)\ \text{algebra-simps})$
have $\text{right-edge-explicit}: ?right\text{-}edge = (\lambda x. (b, g2\ b + x *_R (g1\ b - g2\ b)))$
by $(\text{simp add: twoCisTypeI}(4)\ \text{algebra-simps})$
have $\text{top-edge-explicit}: ?top\text{-}edge = (\lambda x. (a + (b - a) * x, g1\ (a + (b - a) * x)))$
by $(\text{simp add: twoCisTypeI}(4)\ \text{algebra-simps})$
have $\text{left-edge-explicit}: ?left\text{-}edge = (\lambda x. (a, g2\ a + x *_R (g1\ a - g2\ a)))$
by $(\text{simp add: twoCisTypeI}(4)\ \text{algebra-simps})$
show $?thesis$
using $\text{bottom-edge-explicit } \text{right-edge-explicit } \text{top-edge-explicit } \text{left-edge-explicit}$
 twoCisTypeI
by auto
qed

lemma $\text{typeI-twoCube-smooth-edges}$:

assumes $\text{typeI-twoCube } \text{twoC}$

$(k, \gamma) \in \text{boundary } \text{twoC}$

shows $\gamma\ \text{piecewise-C1-differentiable-on } \{0..1\}$

proof –

let $?bottom\text{-}edge = (\lambda x. \text{twoC}(x, 0))$
let $?right\text{-}edge = (\lambda y. \text{twoC}(1, y))$
let $?top\text{-}edge = (\lambda x. \text{twoC}(x, 1))$
let $?left\text{-}edge = (\lambda y. \text{twoC}(0, y))$
obtain $a\ b\ g1\ g2$ **where**
 $\text{twoCisTypeI}: a < b$
 $(\forall x \in \text{cbox } a\ b. g2\ x \leq g1\ x)$
 $\text{cubeImage } \text{twoC} = \{(x,y). x \in \text{cbox } a\ b \wedge y \in \text{cbox } (g2\ x)\ (g1\ x)\}$
 $\text{twoC} = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) * a + x * b)))$
 $g1\ \text{piecewise-C1-differentiable-on } \{a .. b\}$
 $g2\ \text{piecewise-C1-differentiable-on } \{a .. b\}$

```

(λx. twoC(x, 0)) = (λx. (a + (b - a) * x, g2 (a + (b - a) * x)))
(λy. twoC(1, y)) = (λx. (b, g2 b + x *R (g1 b - g2 b)))
(λx. twoC(x, 1)) = (λx. (a + (b - a) * x, g1 (a + (b - a) * x)))
(λy. twoC(0, y)) = (λx. (a, g2 a + x *R (g1 a - g2 a)))
using assms and typeI-cube-explicit-spec[oftwoC]
by auto
have bottom-edge-smooth: (λx. twoC(x, 0)) piecewise-C1-differentiable-on {0..1}
proof -
  have ∀x. (λx. (a + (b - a) * x))-‘ {x} = {(x - a)/(b - a)}
    using twoCisTypeI(1)
    by(auto simp add: Set.vimage-def)
  then have finite-vimg: ∧x. finite({0..1} ∩ (λx. (a + (b - a) * x))-‘ {x}) by
auto
  have scale-shif-smth: (λx. (a + (b - a) * x)) C1-differentiable-on {0..1} using
scale-shift-smooth by auto
  then have scale-shif-pw-smth: (λx. (a + (b - a) * x)) piecewise-C1-differentiable-on
{0..1} using C1-differentiable-imp-piecewise by blast
  have g2-smooth: g2 piecewise-C1-differentiable-on (λx. a + (b - a) * x) ‘ {0..1}
using add-scale-img[OF twoCisTypeI(1)] twoCisTypeI(6) by auto
  have (λx. g2 (a + (b - a) * x)) piecewise-C1-differentiable-on {0..1}
    using piecewise-C1-differentiable-compose[OF scale-shif-pw-smth g2-smooth
finite-vimg]
    by (auto simp add: o-def)
  then have (λx::real. (a + (b - a) * x, g2 (a + (b - a) * x))) piecewise-C1-differentiable-on
{0..1}
    using all-components-smooth-one-pw-smooth-is-pw-smooth[where f = (λx::real.
(a + (b - a) * x, g2 (a + (b - a) * x)))]
    apply (simp only: real-pair-basis)
    by fastforce
  then show ?thesis using twoCisTypeI(7) by auto
qed
have top-edge-smooth: ?top-edge piecewise-C1-differentiable-on {0..1}
proof -
  have ∀x. (λx. (a + (b - a) * x))-‘ {x} = {(x - a)/(b - a)}
    using twoCisTypeI(1)
    by(auto simp add: Set.vimage-def)
  then have finite-vimg: ∧x. finite({0..1} ∩ (λx. (a + (b - a) * x))-‘ {x}) by
auto
  have scale-shif-smth: (λx. (a + (b - a) * x)) C1-differentiable-on {0..1} using
scale-shift-smooth by auto
  then have scale-shif-pw-smth: (λx. (a + (b - a) * x)) piecewise-C1-differentiable-on
{0..1} using C1-differentiable-imp-piecewise by blast
  have g1-smooth: g1 piecewise-C1-differentiable-on (λx. a + (b - a) * x) ‘ {0..1}
using add-scale-img[OF twoCisTypeI(1)] twoCisTypeI(5) by auto
  have (λx. g1 (a + (b - a) * x)) piecewise-C1-differentiable-on {0..1}
    using piecewise-C1-differentiable-compose[OF scale-shif-pw-smth g1-smooth
finite-vimg]
    by (auto simp add: o-def)
  then have (λx. (a + (b - a) * x, g1 (a + (b - a) * x))) piecewise-C1-differentiable-on

```

```

{0..1}
  using all-components-smooth-one-pw-smooth-is-pw-smooth[where  $f = (\lambda x. (a + (b - a) * x, g1 (a + (b - a) * x)))$ ]
  apply (simp only: real-pair-basis)
  by fastforce
  then show ?thesis using twoCisTypeI(9) by auto
qed
have right-edge-smooth: ?right-edge piecewise-C1-differentiable-on {0..1}
proof -
  have  $(\lambda x. (g2 b + x *_R (g1 b - g2 b)))$  C1-differentiable-on {0..1}
  using scale-shift-smooth C1-differentiable-imp-piecewise by auto
  then have  $(\lambda x. (g2 b + x *_R (g1 b - g2 b)))$  piecewise-C1-differentiable-on
{0..1}
  using C1-differentiable-imp-piecewise by fastforce
  then have  $(\lambda x. (b, g2 b + x *_R (g1 b - g2 b)))$  piecewise-C1-differentiable-on
{0..1}
  using pair-prod-smooth-pw-smooth by auto
  then show ?thesis
  using twoCisTypeI(8) by auto
qed
have left-edge-smooth: ?left-edge piecewise-C1-differentiable-on {0..1}
proof -
  have  $(\lambda x. (g2 a + x *_R (g1 a - g2 a)))$  C1-differentiable-on {0..1}
  using scale-shift-smooth by auto
  then have  $(\lambda x. (g2 a + x *_R (g1 a - g2 a)))$  piecewise-C1-differentiable-on
{0..1}
  using C1-differentiable-imp-piecewise by fastforce
  then have  $(\lambda x. (a, g2 a + x *_R (g1 a - g2 a)))$  piecewise-C1-differentiable-on
{0..1}
  using pair-prod-smooth-pw-smooth by auto
  then show ?thesis
  using twoCisTypeI(10) by auto
qed
have  $\gamma = ?bottom-edge \vee \gamma = ?right-edge \vee \gamma = ?top-edge \vee \gamma = ?left-edge$ 
  using assms by (auto simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
  then show ?thesis
  using left-edge-smooth right-edge-smooth top-edge-smooth bottom-edge-smooth
by auto
qed

```

lemma *two-chain-integral-eq-integral-divisible:*

assumes *f-integrable: \forall twoCube \in twoChain. F integrable-on cubeImage twoCube*
and

gen-division: gen-division s (cubeImage ' twoChain) and

valid-two-chain: valid-two-chain twoChain

shows *integral s $F =$ two-chain-integral twoChain F*

proof -

show *integral s $F =$ two-chain-integral twoChain F*

proof (*simp add: two-chain-integral-def*)
have *partial-deriv-integrable*:
 $\forall \text{twoCube} \in \text{twoChain}. ((F \text{ has-integral } (\text{integral } (\text{cubeImage } \text{twoCube}) (F))))$
(*cubeImage twoCube*)
using *f-integrable by auto*
then have *partial-deriv-integrable*:
 $\bigwedge \text{twoCubeImg}. \text{twoCubeImg} \in \text{cubeImage } \text{twoChain} \implies (F \text{ has-integral } (\text{integral } (\text{twoCubeImg}) F)) (\text{twoCubeImg})$
using *Henstock-Kurzweil-Integration.integrable-neg by force*
have *finite-images: finite (cubeImage ' twoChain)*
using *gen-division gen-division-def by auto*
have *negligible-images: pairwise ($\lambda S S'. \text{negligible } (S \cap S')$) (cubeImage ' twoChain)*
using *gen-division by (auto simp add: gen-division-def pairwise-def)*
have *inj: inj-on cubeImage twoChain*
using *valid-two-chain by (simp add: inj-on-def valid-two-chain-def)*
have *integral s F = ($\sum \text{twoCubeImg} \in \text{cubeImage } \text{twoChain}. \text{integral } \text{twoCubeImg } F$)*
using *has-integral-Union[OF finite-images partial-deriv-integrable negligible-images] gen-division*
by (*auto simp add: gen-division-def*)
also have $\dots = (\sum C \in \text{twoChain}. \text{integral } (\text{cubeImage } C) F)$
using *sum.reindex inj by auto*
finally show $\text{integral } s F = (\sum C \in \text{twoChain}. \text{integral } (\text{cubeImage } C) F)$.
qed
qed

definition *only-vertical-division where*
only-vertical-division one-chain two-chain \equiv
 $\exists \mathcal{V} \mathcal{H}. \text{finite } \mathcal{H} \wedge \text{finite } \mathcal{V} \wedge$
 $(\forall (k, \gamma) \in \mathcal{V}. (\exists (k', \gamma') \in \text{two-chain-vertical-boundary two-chain}. (\exists a \in \{0..1\}. \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a \text{ } b \text{ } \gamma' = \gamma))) \wedge$
 $(\text{common-sudiv-exists } (\text{two-chain-horizontal-boundary two-chain}) \mathcal{H})$
 $\vee \text{common-reparam-exists } \mathcal{H} (\text{two-chain-horizontal-boundary two-chain}))$
 \wedge
 $\text{boundary-chain } \mathcal{H} \wedge \text{one-chain} = \mathcal{V} \cup \mathcal{H} \wedge$
 $(\forall (k, \gamma) \in \mathcal{H}. \text{valid-path } \gamma)$

abbreviation *valid-typeI-division s twoChain*
 $\equiv (\forall \text{twoCube} \in \text{twoChain}. \text{typeI-twoCube } \text{twoCube}) \wedge$
 $\text{gen-division } s (\text{cubeImage } \text{twoChain}) \wedge \text{valid-two-chain } \text{twoChain}$

lemma *field-cont-on-typeI-region-cont-on-edges:*
assumes *typeI-twoC: typeI-twoCube twoC*
and *field-cont: continuous-on (cubeImage twoC) F*
and *member-of-boundary: (k, γ) \in boundary twoC*
shows *continuous-on (γ ' {0 .. 1}) F*

proof –

obtain $a\ b\ g1\ g2$ **where**
twoCisTypeI: $a < b$
 $(\forall x \in \text{cbox } a\ b. g2\ x \leq g1\ x)$
cubeImage twoC = $\{(x,y). x \in \text{cbox } a\ b \wedge y \in \text{cbox } (g2\ x)\ (g1\ x)\}$
twoC = $(\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) * a + x * b)))$
g1 piecewise-C1-differentiable-on $\{a .. b\}$
g2 piecewise-C1-differentiable-on $\{a .. b\}$
 $(\lambda x. \text{twoC}(x, 0)) = (\lambda x. (a + (b - a) * x, g2\ (a + (b - a) * x)))$
 $(\lambda y. \text{twoC}(1, y)) = (\lambda x. (b, g2\ b + x *_R (g1\ b - g2\ b)))$
 $(\lambda x. \text{twoC}(x, 1)) = (\lambda x. (a + (b - a) * x, g1\ (a + (b - a) * x)))$
 $(\lambda y. \text{twoC}(0, y)) = (\lambda x. (a, g2\ a + x *_R (g1\ a - g2\ a)))$
using *typeI-twoC* **and** *typeI-cube-explicit-spec[oftwoC]*
by *auto*
let *?bottom-edge* = $(\lambda x. \text{twoC}(x, 0))$
let *?right-edge* = $(\lambda y. \text{twoC}(1, y))$
let *?top-edge* = $(\lambda x. \text{twoC}(x, 1))$
let *?left-edge* = $(\lambda y. \text{twoC}(0, y))$
let *?Dg1* = $\{p. \exists x. x \in \text{cbox } a\ b \wedge p = (x, g1(x))\}$
have *line-is-pair-img*: *?Dg1* = $(\lambda x. (x, g1(x))) \text{ ' } (\text{cbox } a\ b)$
using *image-def* **by** *auto*
have *field-cont-on-top-edge-image*: *continuous-on* *?Dg1* F
by (*rule continuous-on-subset* [*OF field-cont*]) (*auto simp*: *twoCisTypeI(2)*
twoCisTypeI(3))
have *top-edge-is-compos-of-scal-and-g1*:
 $(\lambda x. \text{twoC}(x, 1)) = (\lambda x. (x, g1(x))) \circ (\lambda x. a + (b - a) * x)$
using *twoCisTypeI* **by** *auto*
have *Dg1-is-bot-edge-pathimg*: *path-image* $(\lambda x. \text{twoC}(x, 1)) = ?Dg1$
using *line-is-pair-img* **and** *top-edge-is-compos-of-scal-and-g1* *image-comp* *path-image-def*
add-scale-img **and** *twoCisTypeI(1)*
by (*metis* (*no-types*, *lifting*) *cbox-interval*)
then have *cont-on-top*: *continuous-on* (*path-image* *?top-edge*) F
using *field-cont-on-top-edge-image* **by** *auto*
let *?Dg2* = $\{p. \exists x. x \in \text{cbox } a\ b \wedge p = (x, g2(x))\}$
have *line-is-pair-img*: *?Dg2* = $(\lambda x. (x, g2(x))) \text{ ' } (\text{cbox } a\ b)$
using *image-def* **by** *auto*
have *field-cont-on-bot-edge-image*: *continuous-on* *?Dg2* F
apply (*rule continuous-on-subset* [*OF field-cont*])
using *twoCisTypeI(2)* *twoCisTypeI(3)* **by** *auto*
have *bot-edge-is-compos-of-scal-and-g2*: $(\lambda x. \text{twoC}(x, 0)) = (\lambda x. (x, g2(x))) \circ$
 $(\lambda x. a + (b - a) * x)$
using *twoCisTypeI* **by** *auto*
have *Dg2-is-bot-edge-pathimg*:
path-image $(\lambda x. \text{twoC}(x, 0)) = ?Dg2$
using *line-is-pair-img* **and** *bot-edge-is-compos-of-scal-and-g2* *image-comp* *path-image-def*
add-scale-img **and** *twoCisTypeI(1)*
by (*metis* (*no-types*, *lifting*) *cbox-interval*)
then have *cont-on-bot*: *continuous-on* (*path-image* *?bottom-edge*) F


```

    using field-cont-on-bot-edge-image by auto
  let ?D-left-edge = {p. ∃ y. y ∈ cbox (g2 a) (g1 a) ∧ p = (a, y)}
  have field-cont-on-left-edge-image: continuous-on ?D-left-edge F
    apply (rule continuous-on-subset [OF field-cont])
    using twoCisTypeI(1) twoCisTypeI(3) by auto
  have g2 a ≤ g1 a using twoCisTypeI(1) twoCisTypeI(2) by auto
  then have (λx. g2 a + (g1 a - g2 a) * x) ‘{(0::real)..1} = {g2 a .. g1 a}
    using add-scale-img'[of g2 a g1 a] by blast
  then have left-eq: ?D-left-edge = ?left-edge ‘{0..1}
    unfolding twoCisTypeI(10)
    by (auto simp add: subset-iff image-def set-eq-iff Semiring-Normalization.comm-semiring-1-class.semiring-no)
  then have cont-on-left: continuous-on (path-image ?left-edge) F
    using field-cont-on-left-edge-image path-image-def
    by (metis left-eq field-cont-on-left-edge-image path-image-def)
  let ?D-right-edge = {p. ∃ y. y ∈ cbox (g2 b) (g1 b) ∧ p = (b, y)}
  have field-cont-on-right-edge-image: continuous-on ?D-right-edge F
    apply (rule continuous-on-subset [OF field-cont])
    using twoCisTypeI(1) twoCisTypeI(3) by auto
  have g2 b ≤ g1 b using twoCisTypeI(1) twoCisTypeI(2) by auto
  then have (λx. g2 b + (g1 b - g2 b) * x) ‘{(0::real)..1} = {g2 b .. g1 b}
    using add-scale-img'[of g2 b g1 b] by blast
  then have right-eq: ?D-right-edge = ?right-edge ‘{0..1}
    unfolding twoCisTypeI(8)
    by (auto simp add: subset-iff image-def set-eq-iff Semiring-Normalization.comm-semiring-1-class.semiring-no)
  then have cont-on-right:
    continuous-on (path-image ?right-edge) F
    using field-cont-on-right-edge-image path-image-def
    by (metis right-eq field-cont-on-right-edge-image path-image-def)
  have all-edge-cases:
    (γ = ?bottom-edge ∨ γ = ?right-edge ∨ γ = ?top-edge ∨ γ = ?left-edge)
    using assms by (auto simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
  show ?thesis
    apply (simp add: path-image-def[symmetric])
    using cont-on-top cont-on-bot cont-on-right cont-on-left all-edge-cases
    by blast
qed

```

lemma *typeII-cube-explicit-spec*:

```

  assumes typeII-twoCube twoC
  shows ∃ a b g1 g2. a < b ∧ (∀ x ∈ {a .. b}. g2 x ≤ g1 x) ∧
    cubeImage twoC = {(y, x). x ∈ {a..b} ∧ y ∈ {g2 x .. g1 x}}
    ∧ twoC = (λ(y, x). ((1 - y) * g2 ((1 - x) * a + x * b) + y * g1
      ((1 - x) * a + x * b), (1 - x) * a + x * b))
    ∧ g1 piecewise-C1-differentiable-on {a .. b} ∧ g2 piecewise-C1-differentiable-on
      {a .. b}
    ∧ (λx. twoC(0, x)) = (λx. (g2 (a + (b - a) * x), a + (b - a) * x))
    ∧ (λy. twoC(y, 1)) = (λx. (g2 b + x *R (g1 b - g2 b), b))
    ∧ (λx. twoC(1, x)) = (λx. (g1 (a + (b - a) * x), a + (b - a) * x))

```

$$\wedge (\lambda y. \text{twoC}(y, 0)) = (\lambda x. (g2\ a + x *_R (g1\ a - g2\ a), a))$$

proof –

let *?bottom-edge* = ($\lambda x. \text{twoC}(0, x)$)
let *?right-edge* = ($\lambda y. \text{twoC}(y, 1)$)
let *?top-edge* = ($\lambda x. \text{twoC}(1, x)$)
let *?left-edge* = ($\lambda y. \text{twoC}(y, 0)$)
obtain *a b g1 g2* **where**
twoCisTypeII: $a < b$
 $(\forall x \in \text{cbox } a\ b. g2\ x \leq g1\ x)$
cubeImage twoC = $\{(y, x). x \in \text{cbox } a\ b \wedge y \in \text{cbox } (g2\ x)\ (g1\ x)\}$
twoC = ($\lambda(y, x). ((1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) * a + x * b), (1 - x) * a + x * b)$)
g1 piecewise-C1-differentiable-on $\{a .. b\}$
g2 piecewise-C1-differentiable-on $\{a .. b\}$
using *assms and typeII-twoCubeImg[oftwoC]* **by** *auto*
have *bottom-edge-explicit*: *?bottom-edge* = ($\lambda x. (g2\ (a + (b - a) * x), a + (b - a) * x)$)
by (*simp add: twoCisTypeII(4) algebra-simps*)
have *right-edge-explicit*: *?right-edge* = ($\lambda x. (g2\ b + x *_R (g1\ b - g2\ b), b)$)
by (*simp add: twoCisTypeII(4) algebra-simps*)
have *top-edge-explicit*: *?top-edge* = ($\lambda x. (g1\ (a + (b - a) * x), a + (b - a) * x)$)
by (*simp add: twoCisTypeII(4) algebra-simps*)
have *left-edge-explicit*: *?left-edge* = ($\lambda x. (g2\ a + x *_R (g1\ a - g2\ a), a)$)
by (*simp add: twoCisTypeII(4) algebra-simps*)
show *?thesis*
using *bottom-edge-explicit right-edge-explicit top-edge-explicit left-edge-explicit twoCisTypeII*
by *auto*
qed

lemma *typeII-twoCube-smooth-edges*:

assumes *typeII-twoCube twoC* $(k, \gamma) \in \text{boundary twoC}$

shows γ *piecewise-C1-differentiable-on* $\{0..1\}$

proof –

let *?bottom-edge* = ($\lambda x. \text{twoC}(0, x)$)
let *?right-edge* = ($\lambda y. \text{twoC}(y, 1)$)
let *?top-edge* = ($\lambda x. \text{twoC}(1, x)$)
let *?left-edge* = ($\lambda y. \text{twoC}(y, 0)$)
obtain *a b g1 g2* **where**
twoCisTypeII: $a < b$
 $(\forall x \in \text{cbox } a\ b. g2\ x \leq g1\ x)$
cubeImage twoC = $\{(y, x). x \in \text{cbox } a\ b \wedge y \in \text{cbox } (g2\ x)\ (g1\ x)\}$
twoC = ($\lambda(y, x). ((1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) * a + x * b), (1 - x) * a + x * b)$)
g1 piecewise-C1-differentiable-on $\{a .. b\}$
g2 piecewise-C1-differentiable-on $\{a .. b\}$
 $(\lambda x. \text{twoC}(0, x)) = (\lambda x. (g2\ (a + (b - a) * x), a + (b - a) * x))$
 $(\lambda y. \text{twoC}(y, 1)) = (\lambda x. (g2\ b + x *_R (g1\ b - g2\ b), b))$

```

    (λx. twoC(1, x)) = (λx. (g1 (a + (b - a) * x), a + (b - a) * x))
    (λy. twoC(y, 0)) = (λx. (g2 a + x *R (g1 a - g2 a), a))
    using assms and typeII-cube-explicit-spec[oftwoC]
    by auto
  have bottom-edge-smooth: ?bottom-edge piecewise-C1-differentiable-on {0..1}
  proof -
    have ∀x. (λx. (a + (b - a) * x)) -' {x} = {(x - a)/(b - a)}
      using twoCisTypeII(1) by auto
    then have finite-vimg: ∧x. finite({0..1} ∩ (λx. (a + (b - a) * x)) -' {x}) by
    auto
    have scale-shif-pw-smth: (λx. (a + (b - a) * x)) piecewise-C1-differentiable-on
    {0..1}
      using scale-shift-smooth C1-differentiable-imp-piecewise by blast
    have g2-smooth: g2 piecewise-C1-differentiable-on (λx. a + (b - a) * x) ' {0..1}
  using add-scale-img[OF twoCisTypeII(1)] twoCisTypeII(6) by auto
    have (λx. g2 (a + (b - a) * x)) piecewise-C1-differentiable-on {0..1}
      using piecewise-C1-differentiable-compose[OF scale-shif-pw-smth g2-smooth
    finite-vimg]
      by (auto simp add: o-def)
    then have (λx::real. (g2 (a + (b - a) * x), a + (b - a) * x)) piece-
    wise-C1-differentiable-on {0..1}
      using all-components-smooth-one-pw-smooth-is-pw-smooth[where f = (λx::real.
    (g2 (a + (b - a) * x), a + (b - a) * x))]
      by (fastforce simp add: real-pair-basis)
    then show ?thesis using twoCisTypeII(7) by auto
  qed
  have top-edge-smooth: ?top-edge piecewise-C1-differentiable-on {0..1}
  proof -
    have ∀x. (λx. (a + (b - a) * x)) -' {x} = {(x - a)/(b - a)}
      using twoCisTypeII(1) by auto
    then have finite-vimg: ∧x. finite({0..1} ∩ (λx. (a + (b - a) * x)) -' {x}) by
    auto
    have scale-shif-pw-smth: (λx. (a + (b - a) * x)) piecewise-C1-differentiable-on
    {0..1}
      using scale-shift-smooth C1-differentiable-imp-piecewise by blast
    have g1-smooth: g1 piecewise-C1-differentiable-on (λx. a + (b - a) * x) ' {0..1}
  using add-scale-img[OF twoCisTypeII(1)] twoCisTypeII(5) by auto
    have (λx. g1 (a + (b - a) * x)) piecewise-C1-differentiable-on {0..1}
      using piecewise-C1-differentiable-compose[OF scale-shif-pw-smth g1-smooth
    finite-vimg]
      by (auto simp add: o-def)
    then have (λx::real. (g1 (a + (b - a) * x), a + (b - a) * x)) piece-
    wise-C1-differentiable-on {0..1}
      using all-components-smooth-one-pw-smooth-is-pw-smooth[where f = (λx::real.
    (g1 (a + (b - a) * x), a + (b - a) * x))]
      by (fastforce simp add: real-pair-basis)
    then show ?thesis using twoCisTypeII(9) by auto
  qed
  have right-edge-smooth: ?right-edge piecewise-C1-differentiable-on {0..1}

```

proof –
have $(\lambda x. (g2\ b + x *_R (g1\ b - g2\ b)))$ *piecewise-C1-differentiable-on* $\{0..1\}$
by *(simp add: C1-differentiable-imp-piecewise)*
then have $(\lambda x. (g2\ b + x *_R (g1\ b - g2\ b), b))$ *piecewise-C1-differentiable-on*
 $\{0..1\}$
using *all-components-smooth-one-pw-smooth-is-pw-smooth*[of $(1,0)$ $(\lambda x. (g2\ b + x *_R (g1\ b - g2\ b), b))$]
by *(auto simp add: real-pair-basis)*
then show *?thesis*
using *twoCisTypeII(8)* **by** *auto*
qed
have *left-edge-smooth: ?left-edge piecewise-C1-differentiable-on* $\{0..1\}$
proof –
have $0: (\lambda x. (g2\ a + x *_R (g1\ a - g2\ a)))$ *C1-differentiable-on* $\{0..1\}$
using *C1-differentiable-imp-piecewise* **by** *fastforce*
have $(\lambda x. (g2\ a + x *_R (g1\ a - g2\ a), a))$ *piecewise-C1-differentiable-on* $\{0..1\}$
using *C1-differentiable-imp-piecewise*[OF *C1-differentiable-on-pair*[OF 0 *C1-differentiable-on-const*[of
 a $\{0..1\}$]]]
by *force*
then show *?thesis*
using *twoCisTypeII(10)* **by** *auto*
qed
have $\gamma = ?bottom-edge \vee \gamma = ?right-edge \vee \gamma = ?top-edge \vee \gamma = ?left-edge$
using *assms* **by** *(auto simp add: boundary-def horizontal-boundary-def verti-
cal-boundary-def)*
then show *?thesis*
using *left-edge-smooth right-edge-smooth top-edge-smooth bottom-edge-smooth*
by *auto*
qed

lemma *field-cont-on-typeII-region-cont-on-edges:*

assumes *typeII-twoC:*
typeII-twoCube twoC **and**
field-cont:
continuous-on (cubeImage twoC) F **and**
member-of-boundary:
 $(k,\gamma) \in \text{boundary } twoC$
shows *continuous-on* $(\gamma \text{ ‘ } \{0 .. 1\}) F$
proof –
obtain $a\ b\ g1\ g2$ **where**
twoCisTypeII: a < b
 $(\forall x \in \text{cbox } a\ b. g2\ x \leq g1\ x)$
 $\text{cubeImage } twoC = \{(y, x). x \in \text{cbox } a\ b \wedge y \in \text{cbox } (g2\ x)\ (g1\ x)\}$
 $twoC = (\lambda(y, x). ((1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) * a$
 $+ x * b), (1 - x) * a + x * b))$
 $g1$ *piecewise-C1-differentiable-on* $\{a .. b\}$
 $g2$ *piecewise-C1-differentiable-on* $\{a .. b\}$
 $(\lambda x. twoC(0, x)) = (\lambda x. (g2\ (a + (b - a) * x), a + (b - a) * x))$
 $(\lambda y. twoC(y, 1)) = (\lambda x. (g2\ b + x *_R (g1\ b - g2\ b), b))$

```

    (λx. twoC(1, x)) = (λx. (g1 (a + (b - a) * x), a + (b - a) * x))
    (λy. twoC(y, 0)) = (λx. (g2 a + x *R (g1 a - g2 a), a))
    using typeII-twoC and typeII-cube-explicit-spec[oftwoC]
    by auto
  let ?bottom-edge = (λx. twoC(0, x))
  let ?right-edge = (λy. twoC(y, 1))
  let ?top-edge = (λx. twoC(1, x))
  let ?left-edge = (λy. twoC(y, 0))
  let ?Dg1 = {p. ∃x. x ∈ cbox a b ∧ p = (g1(x), x)}
  have line-is-pair-img: ?Dg1 = (λx. (g1(x), x)) ‘ (cbox a b)
    using image-def by auto
  have field-cont-on-top-edge-image: continuous-on ?Dg1 F
    by (rule continuous-on-subset [OF field-cont]) (auto simp: twoCisTypeII(2)
twoCisTypeII(3))
  have top-edge-is-compos-of-scal-and-g1:
    (λx. twoC(1, x)) = (λx. (g1(x), x)) ∘ (λx. a + (b - a) * x)
    using twoCisTypeII by auto
  have Dg1-is-bot-edge-pathimg:
    path-image (λx. twoC(1, x)) = ?Dg1
  using line-is-pair-img and top-edge-is-compos-of-scal-and-g1 image-comp path-image-def
add-scale-img and twoCisTypeII(1)
  by (metis (no-types, lifting) cbox-interval)
  then have cont-on-top: continuous-on (path-image ?top-edge) F
    using field-cont-on-top-edge-image by auto
  let ?Dg2 = {p. ∃x. x ∈ cbox a b ∧ p = (g2(x), x)}
  have line-is-pair-img: ?Dg2 = (λx. (g2(x), x)) ‘ (cbox a b)
    using image-def by auto
  have field-cont-on-bot-edge-image: continuous-on ?Dg2 F
    by (rule continuous-on-subset [OF field-cont]) (auto simp add: twoCisTypeII(2)
twoCisTypeII(3))
  have bot-edge-is-compos-of-scal-and-g2:
    (λx. twoC(0, x)) = (λx. (g2(x), x)) ∘ (λx. a + (b - a) * x)
    using twoCisTypeII by auto
  have Dg2-is-bot-edge-pathimg: path-image (λx. twoC(0, x)) = ?Dg2
  unfolding path-image-def
  using line-is-pair-img and bot-edge-is-compos-of-scal-and-g2 image-comp add-scale-img
[OF ‹a < b›]
  by (metis (no-types, lifting) box-real(2))
  then have cont-on-bot: continuous-on (path-image ?bottom-edge) F
    using field-cont-on-bot-edge-image
    by auto
  let ?D-left-edge = {p. ∃y. y ∈ cbox (g2 a) (g1 a) ∧ p = (y, a)}
  have field-cont-on-left-edge-image: continuous-on ?D-left-edge F
    apply (rule continuous-on-subset [OF field-cont])
    using twoCisTypeII(1) twoCisTypeII(3) by auto
  have g2 a ≤ g1 a using twoCisTypeII(1) twoCisTypeII(2) by auto
  then have (λx. g2 a + x * (g1 a - g2 a)) ‘ {(0::real)..1} = {g2 a .. g1 a}
    using add-scale-img'[of g2 a g1 a] by (auto simp add: ac-simps)
  with ‹g2 a ≤ g1 a› have left-eq: ?D-left-edge = ?left-edge ‘ {0..1}

```

```

  by (simp only: twoCisTypeII(10)) auto
then have cont-on-left: continuous-on (path-image ?left-edge) F
  using field-cont-on-left-edge-image path-image-def
  by (metis left-eq field-cont-on-left-edge-image path-image-def)
let ?D-right-edge = {p.  $\exists y. y \in \text{cbox } (g2 \ b) \ (g1 \ b) \wedge p = (y, \ b)$ }
have field-cont-on-left-edge-image: continuous-on ?D-right-edge F
  apply (rule continuous-on-subset [OF field-cont])
  using twoCisTypeII(1) twoCisTypeII(3) by auto
have  $g2 \ b \leq g1 \ b$  using twoCisTypeII(1) twoCisTypeII(2) by auto
then have  $(\lambda x. g2 \ b + x * (g1 \ b - g2 \ b)) \ ' \ \{(0::\text{real})..1\} = \{g2 \ b .. g1 \ b\}$ 
  using add-scale-imag'[of  $g2 \ b \ g1 \ b$ ] by (auto simp add: ac-simps)
with  $\langle g2 \ b \leq g1 \ b \rangle$  have right-eq: ?D-right-edge = ?right-edge '  $\{0..1\}$ 
  by (simp only: twoCisTypeII(8)) auto
then have cont-on-right:
  continuous-on (path-image ?right-edge) F
  using field-cont-on-left-edge-image path-image-def
  by (metis right-eq field-cont-on-left-edge-image path-image-def)
have all-edge-cases:
   $(\gamma = ?\text{bottom-edge} \vee \gamma = ?\text{right-edge} \vee \gamma = ?\text{top-edge} \vee \gamma = ?\text{left-edge})$ 
  using assms unfolding boundary-def horizontal-boundary-def vertical-boundary-def
by blast
show ?thesis
  apply (simp add: path-image-def[symmetric])
  using cont-on-top cont-on-bot cont-on-right cont-on-left all-edge-cases
  by blast
qed

```

lemma *two-cube-boundary-is-boundary*: boundary-chain (boundary C)
 by (auto simp add: boundary-chain-def boundary-def horizontal-boundary-def vertical-boundary-def)

lemma *common-boundary-subdiv-exists-refl*:
 assumes $\forall (k,\gamma) \in \text{boundary } \text{two}C. \text{valid-path } \gamma$
 shows *common-boundary-sudivision-exists* (boundary twoC) (boundary twoC)
 using assms chain-subdiv-chain-refl common-boundary-sudivision-exists-def two-cube-boundary-is-boundary
 by blast

lemma *common-boundary-subdiv-exists-refl'*:
 assumes $\forall (k,\gamma) \in C. \text{valid-path } \gamma$
 boundary-chain (C::($\text{int} \times (\text{real} \Rightarrow \text{real} \times \text{real})$)) set)
 shows *common-boundary-sudivision-exists* (C) (C)
 using assms chain-subdiv-chain-refl common-boundary-sudivision-exists-def by
 blast

lemma *gen-common-boundary-subdiv-exists-refl-twochain-boundary*:
 assumes $\forall (k,\gamma) \in C. \text{valid-path } \gamma$
 boundary-chain (C::($\text{int} \times (\text{real} \Rightarrow \text{real} \times \text{real})$)) set)
 shows *common-sudiv-exists* (C) (C)
 using assms chain-subdiv-chain-refl common-boundary-sudivision-exists-def com-

mon-subdiv-imp-gen-common-subdiv **by** *blast*

lemma *two-chain-boundary-is-boundary-chain*:

shows *boundary-chain* (*two-chain-boundary twoChain*)

by (*simp add: boundary-chain-def two-chain-boundary-def boundary-def horizontal-boundary-def vertical-boundary-def*)

lemma *typeI-edges-are-valid-paths*:

assumes *typeI-twoCube twoC* $(k, \gamma) \in \text{boundary twoC}$

shows *valid-path* γ

using *typeI-twoCube-smooth-edges[OF assms]* *C1-differentiable-imp-piecewise*

by (*auto simp: valid-path-def*)

lemma *typeII-edges-are-valid-paths*:

assumes *typeII-twoCube twoC* $(k, \gamma) \in \text{boundary twoC}$

shows *valid-path* γ

using *typeII-twoCube-smooth-edges[OF assms]* *C1-differentiable-imp-piecewise*

by (*auto simp: valid-path-def*)

lemma *finite-two-chain-vertical-boundary*:

assumes *finite two-chain*

shows *finite* (*two-chain-vertical-boundary two-chain*)

using *assms* **by** (*simp add: two-chain-vertical-boundary-def vertical-boundary-def*)

lemma *finite-two-chain-horizontal-boundary*:

assumes *finite two-chain*

shows *finite* (*two-chain-horizontal-boundary two-chain*)

using *assms* **by** (*simp add: two-chain-horizontal-boundary-def horizontal-boundary-def*)

locale *R2* =

fixes *i j*

assumes *i-is-x-axis*: $i = (1::\text{real}, 0::\text{real})$ **and**

j-is-y-axis: $j = (0::\text{real}, 1::\text{real})$

begin

lemma *analytically-valid-y*:

assumes *analytically-valid s F i*

shows $(\lambda x. \text{integral UNIV } (\lambda y. (\text{partial-vector-derivative } F \ i) \ (y, x) * (\text{indicator } s \ (y, x)))) \in \text{borel-measurable lborel}$

proof –

have $\{(1::\text{real}, 0::\text{real}), (0, 1)\} - \{(1, 0)\} = \{(0::\text{real}, 1::\text{real})\}$

by *force*

with *assms* **show** *?thesis*

using *assms* **by** (*simp add: real-pair-basis analytically-valid-def i-is-x-axis*

j-is-y-axis)

qed

lemma *analytically-valid-x*:

assumes *analytically-valid s F j*

shows $(\lambda x. \text{integral UNIV } (\lambda y. ((\text{partial-vector-derivative } F \ j) \ (x, y)) * (\text{indicator } s \ (x, y)))) \in \text{borel-measurable lborel}$

proof –

have $\{(1::\text{real}, 0::\text{real}), (0, 1)\} - \{(0, 1)\} = \{(1::\text{real}, 0::\text{real})\}$

by force

with *assms* **show** *?thesis*

by (*simp add: real-pair-basis analytically-valid-def i-is-x-axis j-is-y-axis*)

qed

lemma *Greens-thm-type-I:*

fixes $F:: ((\text{real} * \text{real}) \Rightarrow (\text{real} * \text{real}))$ **and**

$\text{gamma1 } \text{gamma2 } \text{gamma3 } \text{gamma4} :: (\text{real} \Rightarrow (\text{real} * \text{real}))$ **and**

$a:: \text{real}$ **and** $b:: \text{real}$ **and**

$g1:: (\text{real} \Rightarrow \text{real})$ **and** $g2:: (\text{real} \Rightarrow \text{real})$

assumes *Dy-def: Dy-pair* = $\{(x::\text{real}, y) . x \in \text{cbox } a \ b \wedge y \in \text{cbox } (g2 \ x) \ (g1 \ x)\}$

and

gamma1-def: gamma1 = $(\lambda x. (a + (b - a) * x, g2(a + (b - a) * x)))$ **and**

gamma1-smooth: gamma1 *piecewise-C1-differentiable-on* $\{0..1\}$ **and**

gamma2-def: gamma2 = $(\lambda x. (b, g2(b) + x *_{\mathbb{R}} (g1(b) - g2(b))))$ **and**

gamma3-def: gamma3 = $(\lambda x. (a + (b - a) * x, g1(a + (b - a) * x)))$ **and**

gamma3-smooth: gamma3 *piecewise-C1-differentiable-on* $\{0..1\}$ **and**

gamma4-def: gamma4 = $(\lambda x. (a, g2(a) + x *_{\mathbb{R}} (g1(a) - g2(a))))$ **and**

F-i-analytically-valid: analytically-valid Dy-pair $(\lambda p. F(p) \cdot i)$ *j* **and**

g2-leq-g1: $\forall x \in \text{cbox } a \ b. (g2 \ x) \leq (g1 \ x)$ **and**

a-lt-b: $a < b$

shows $(\text{line-integral } F \ \{i\} \ \text{gamma1}) +$

$(\text{line-integral } F \ \{i\} \ \text{gamma2}) -$

$(\text{line-integral } F \ \{i\} \ \text{gamma3}) -$

$(\text{line-integral } F \ \{i\} \ \text{gamma4})$

$= (\text{integral } \text{Dy-pair} \ (\lambda a. - (\text{partial-vector-derivative } (\lambda p. F(p) \cdot i) \ j$

$a)))$

line-integral-exists $F \ \{i\} \ \text{gamma4}$

line-integral-exists $F \ \{i\} \ \text{gamma3}$

line-integral-exists $F \ \{i\} \ \text{gamma2}$

line-integral-exists $F \ \{i\} \ \text{gamma1}$

proof –

let *?F-b'* = *partial-vector-derivative* $(\lambda a. (F \ a) \cdot i) \ j$

have *F-first-is-continuous: continuous-on Dy-pair* $(\lambda a. F(a) \cdot i)$

using *F-i-analytically-valid*

by (*auto simp add: analytically-valid-def*)

let *?f* = $(\lambda x. \text{if } x \in (\text{Dy-pair}) \text{ then } (\text{partial-vector-derivative } (\lambda a. (F \ a) \cdot i) \ j) \ x \ \text{else } 0)$

have *f-lesbegue-integrable: integrable lborel ?f*

using *F-i-analytically-valid*

by (*auto simp add: analytically-valid-def indic-ident*)

have *partially-vec-diff: $\forall a \in \text{Dy-pair}. \text{partially-vector-differentiable } (\lambda a. (F \ a) \cdot i) \ j \ a$*

using *F-i-analytically-valid*

by (*auto simp add: analytically-valid-def indicator-def*)

have *x-axis-integral-measurable*: $(\lambda x. \text{integral UNIV } (\lambda y. ?f(x, y))) \in \text{borel-measurable lborel}$
proof –
have $(\lambda p. (?F-b' p) * \text{indicator } (Dy\text{-pair}) p) = (\lambda x. \text{if } x \in (Dy\text{-pair}) \text{ then } (?F-b') x \text{ else } 0)$
using *indic-ident*[of *?F-b'*] **by** *auto*
then have $\bigwedge x y. ?F-b'(x, y) * \text{indicator } (Dy\text{-pair}) (x, y) = (\lambda x. \text{if } x \in (Dy\text{-pair}) \text{ then } (?F-b') x \text{ else } 0) (x, y)$
then *show* *?thesis*
using *analytically-valid-x*[*OF F-i-analytically-valid*]
by (*auto simp add: indicator-def*)
qed
have *F-partially-differentiable*: $\forall a \in Dy\text{-pair}. \text{has-partial-vector-derivative } (\lambda x. (F x) \cdot i) j (?F-b' a) a$
using *partial-vector-derivative-works-partially-vec-diff*
by *fastforce*
have *g1-g2-continuous*: *continuous-on* (*cbox a b*) *g1*
continuous-on (*cbox a b*) *g2*
proof –
have *shift-scale-cont*: *continuous-on* $\{a..b\}$ $(\lambda x. (x - a) * (1 / (b - a)))$
by (*intro continuous-intros*)
have *shift-scale-inv*: $(\lambda x. a + (b - a) * x) \circ (\lambda x. (x - a) * (1 / (b - a))) = \text{id}$
using *a-lt-b* **by** (*auto simp add: o-def*)
have *img-shift-scale*: $(\lambda x. (x - a) * (1 / (b - a))) \{a..b\} = \{0..1\}$
using *a-lt-b* **apply** (*auto simp: divide-simps image-iff*)
apply (*rule-tac x=x * (b - a) + a in bexI*)
using *le-diff-eq* **by** *fastforce+*
have *gamma1-y-component*: $(\lambda x. g2(a + (b - a) * x)) = g2 \circ (\lambda x. (a + (b - a) * x))$
by *auto*
have *continuous-on* $\{0..1\}$ $(\lambda x. g2(a + (b - a) * x))$
using *continuous-on-inner*[*OF piecewise-C1-differentiable-on-imp-continuous-on* [*OF gamma1-smooth*], of $(\lambda x. j)$, *OF continuous-on-const*]
by (*simp add: gamma1-def j-is-y-axis*)
then have *continuous-on* $\{a..b\}$ $((\lambda x. g2(a + (b - a) * x)) \circ (\lambda x. (x - a) * (1 / (b - a))))$
using *img-shift-scale continuous-on-compose shift-scale-cont*
by *force*
then have *continuous-on* $\{a..b\}$ $(g2 \circ (\lambda x. (a + (b - a) * x)) \circ (\lambda x. (x - a) * (1 / (b - a))))$
using *gamma1-y-component* **by** *auto*
then show *continuous-on* (*cbox a b*) *g2*
using *a-lt-b* **by** (*simp add: shift-scale-inv*)
have *gamma3-y-component*: $(\lambda x. g1(a + (b - a) * x)) = g1 \circ (\lambda x. (a + (b - a) * x))$
by *auto*
have *continuous-on* $\{0..1\}$ $(\lambda x. g1(a + (b - a) * x))$
using *continuous-on-inner*[*OF piecewise-C1-differentiable-on-imp-continuous-on* [*OF*

```

gamma3-smooth], of (λx. j), OF continuous-on-const]
  by (simp add: gamma3-def j-is-y-axis)
  then have continuous-on {a..b} ((λx. g1(a + (b - a) * x)) ∘ (λx. (x -
a)*(1/(b-a))))
  using img-shift-scale continuous-on-compose shift-scale-cont
  by force
  then have continuous-on {a..b} (g1 ∘ (λx.(a + (b - a) * x)) ∘ (λx. (x -
a)*(1/(b-a))))
  using gamma3-y-component by auto
  then show continuous-on (cbox a b) g1
  using a-lt-b by (simp add: shift-scale-inv)
qed
have g2-scale-j-contin: continuous-on (cbox a b) (λx. (0, g2 x))
  by (intro continuous-intros g1-g2-continuous)
let ?Dg2 = {p. ∃x. x ∈ cbox a b ∧ p = (x, g2(x))}
have line-is-pair-img: ?Dg2 = (λx. (x, g2(x))) ‘ (cbox a b)
  using image-def by auto
have g2-path-continuous: continuous-on (cbox a b) (λx. (x, g2(x)))
  by (intro continuous-intros g1-g2-continuous)
have field-cont-on-gamma1-image: continuous-on ?Dg2 (λa. F(a) · i)
  apply (rule continuous-on-subset [OF F-first-is-continuous])
  by (auto simp add: Dy-def g2-leq-g1)
have gamma1-is-compos-of-scal-and-g2:
  gamma1 = (λx. (x, g2(x))) ∘ (λx. a + (b - a) * x)
  using gamma1-def by auto
have add-scale-img:
  (λx. a + (b - a) * x) ‘ {0 .. 1} = {a .. b} using add-scale-img and a-lt-b by
auto
  then have Dg2-is-gamma1-pathimg: path-image gamma1 = ?Dg2
  by (metis (no-types, lifting) box-real(2) gamma1-is-compos-of-scal-and-g2 im-
age-comp line-is-pair-img path-image-def)
  have Base-vecs: i ∈ Basis j ∈ Basis i ≠ j
  using real-pair-basis and i-is-x-axis and j-is-y-axis by auto
  have gamma1-as-euclid-space-fun: gamma1 = (λx. (a + (b - a) * x) *R i + (0,
g2 (a + (b - a) * x)))
  using i-is-x-axis gamma1-def by auto
  have 0: line-integral F {i} gamma1 = integral (cbox a b) (λx. F(x, g2(x)) · i)
  line-integral-exists F {i} gamma1
  using line-integral-on-pair-path-strong [OF norm-Basis[OF Base-vecs(1)] -
gamma1-as-euclid-space-fun, of F]
  gamma1-def gamma1-smooth g2-scale-j-contin a-lt-b add-scale-img
  Dg2-is-gamma1-pathimg and field-cont-on-gamma1-image
  by (auto simp: pathstart-def pathfinish-def i-is-x-axis)
  then show (line-integral-exists F {i} gamma1) by metis
  have gamma2-x-const: ∀x. gamma2 x · i = b
  by (simp add: i-is-x-axis gamma2-def)
  have 1: (line-integral F {i} gamma2) = 0 (line-integral-exists F {i} gamma2)
  using line-integral-on-pair-straight-path[OF gamma2-x-const] straight-path-differentiable-x
gamma2-def

```

by (auto simp add: mult.commute)
 then show (line-integral-exists $F \{i\}$ gamma2) by metis
 have continuous-on (cbox a b) ($\lambda x. F(x, g2(x)) \cdot i$)
 using line-is-pair-img and g2-path-continuous and field-cont-on-gamma1-image
 Topological-Spaces.continuous-on-compose i-is-x-axis j-is-y-axis
 by auto
 then have 6: ($\lambda x. F(x, g2(x)) \cdot i$) integrable-on (cbox a b)
 using integrable-continuous [of a b ($\lambda x. F(x, g2(x)) \cdot i$)] by auto
 have g1-scale-j-contin: continuous-on (cbox a b) ($\lambda x. (0, g1 x)$)
 by (intro continuous-intros g1-g2-continuous)
 let ?Dg1 = { $p. \exists x. x \in \text{cbox } a \text{ } b \wedge p = (x, g1(x))$ }
 have line-is-pair-img: ?Dg1 = ($\lambda x. (x, g1(x))$) ‘ (cbox a b)
 using image-def by auto
 have g1-path-continuous: continuous-on (cbox a b) ($\lambda x. (x, g1(x))$)
 by (intro continuous-intros g1-g2-continuous)
 have field-cont-on-gamma3-image: continuous-on ?Dg1 ($\lambda a. F(a) \cdot i$)
 apply (rule continuous-on-subset [OF F-first-is-continuous])
 by (auto simp add: Dy-def g2-leq-g1)
 have gamma3-is-compos-of-scal-and-g1:
 gamma3 = ($\lambda x. (x, g1(x))$) \circ ($\lambda x. a + (b - a) * x$)
 using gamma3-def by auto
 then have Dg1-is-gamma3-pathimg: path-image gamma3 = ?Dg1
 by (metis (no-types, lifting) box-real(2) image-comp line-is-pair-img local.add-scale-img
 path-image-def)
 have Base-vecs: $i \in \text{Basis } j \in \text{Basis } i \neq j$
 using real-pair-basis and i-is-x-axis and j-is-y-axis by auto
 have gamma3-as-euclid-space-fun: gamma3 = ($\lambda x. (a + (b - a) * x) *_R i + (0,$
 $g1 (a + (b - a) * x))$)
 using i-is-x-axis gamma3-def by auto
 have 2: line-integral $F \{i\}$ gamma3 = integral (cbox a b) ($\lambda x. F(x, g1(x)) \cdot i$)
 line-integral-exists $F \{i\}$ gamma3
 using line-integral-on-pair-path-strong [OF norm-Basis[OF Base-vecs(1)] -
 gamma3-as-euclid-space-fun, of F]
 gamma3-def and gamma3-smooth and g1-scale-j-contin and a-lt-b and
 add-scale-img
 Dg1-is-gamma3-pathimg and field-cont-on-gamma3-image
 by (auto simp: pathstart-def pathfinish-def i-is-x-axis)
 then show (line-integral-exists $F \{i\}$ gamma3) by metis
 have gamma4-x-const: $\forall x. \text{gamma4 } x \cdot i = a$
 using gamma4-def
 by (auto simp add: real-inner-class.inner-add-left inner-not-same-Basis i-is-x-axis)
 have 3: (line-integral $F \{i\}$ gamma4) = 0 (line-integral-exists $F \{i\}$ gamma4)
 using line-integral-on-pair-straight-path[OF gamma4-x-const] straight-path-differentiable-x
 gamma4-def
 by (auto simp add: mult.commute)
 then show (line-integral-exists $F \{i\}$ gamma4)
 by metis
 have continuous-on (cbox a b) ($\lambda x. F(x, g1(x)) \cdot i$)
 using line-is-pair-img and g1-path-continuous and field-cont-on-gamma3-image

continuous-on-compose i-is-x-axis j-is-y-axis
by auto
then have 7: $(\lambda x. F(x, g1(x)) \cdot i)$ *integrable-on* $(cbox\ a\ b)$
using *integrable-continuous* [of a b $(\lambda x. F(x, g1(x)) \cdot i)$]
by auto
have *partial-deriv-one-d-integrable*:
 $((\lambda y. ?F-b'(xc, y)) \text{ has-integral } F(xc, g1(xc)) \cdot i - F(xc, g2(xc)) \cdot i) (cbox\ (g2\ xc)\ (g1\ xc))$
if $xc \in cbox\ a\ b$ **for** xc
proof –
have $\{(xc', y). y \in cbox\ (g2\ xc)\ (g1\ xc) \wedge xc' = xc\} \subseteq Dy\text{-pair}$
using *that by* $(auto\ simp\ add: Dy\text{-def})$
then show $((\lambda y. ?F-b'(xc, y)) \text{ has-integral } F(xc, g1\ xc) \cdot i - F(xc, g2\ xc) \cdot i) (cbox\ (g2\ xc)\ (g1\ xc))$
using *that and Base-vecs and F-partially-differentiable and Dy-def [symmetric] and g2-leq-g1 and fundamental-theorem-of-calculus-partial-vector [of g2 xc g1 xc j i xc *_R i Dy-pair F ?F-b']*
by $(auto\ simp\ add: Groups.ab\text{-semigroup-add-class.add.commute}\ i\text{-is-x-axis}\ j\text{-is-y-axis})$
qed
have *partial-deriv-integrable*: $(?F-b')$ *integrable-on* $Dy\text{-pair}$
by $(simp\ add: F\text{-i-analytically-valid}\ \text{analytically-valid-imp-part-deriv-integrable-on})$
have 4: *integral Dy-pair ?F-b'*
 $= \text{integral } (cbox\ a\ b) (\lambda x. \text{integral } (cbox\ (g2\ x)\ (g1\ x)) (\lambda y. ?F-b'(x, y)))$
proof –
have *x-axis-gauge-integrable*:
 $\bigwedge x. (\lambda y. ?f(x, y)) \text{ integrable-on } UNIV$
proof –
fix $x::real$
have $\forall x. x \notin cbox\ a\ b \longrightarrow (\lambda y. ?f(x, y)) = (\lambda y. 0)$
by $(auto\ simp\ add: Dy\text{-def})$
then have *f-integrable-x-not-in-range*:
 $\forall x. x \notin cbox\ a\ b \longrightarrow (\lambda y. ?f(x, y)) \text{ integrable-on } UNIV$
by $(simp\ add: integrable-0)$
let $?F-b'\text{-oneD} = (\lambda x. (\lambda y. \text{if } y \in (cbox\ (g2\ x)\ (g1\ x)) \text{ then } ?F-b'(x, y) \text{ else } 0))$
have *f-value-x-in-range*: $\forall x \in cbox\ a\ b. ?F-b'\text{-oneD } x = (\lambda y. ?f(x, y))$
by $(auto\ simp\ add: Dy\text{-def})$
have $\forall x \in cbox\ a\ b. ?F-b'\text{-oneD } x \text{ integrable-on } UNIV$
using *has-integral-integrable integrable-restrict-UNIV partial-deriv-one-d-integrable*
by *blast*
then have *f-integrable-x-in-range*:
 $\forall x. x \in cbox\ a\ b \longrightarrow (\lambda y. ?f(x, y)) \text{ integrable-on } UNIV$
using *f-value-x-in-range by auto*
show $(\lambda y. ?f(x, y)) \text{ integrable-on } UNIV$
using *f-integrable-x-not-in-range and f-integrable-x-in-range by auto*
qed
have *arg*: $(\lambda a. \text{if } a \in Dy\text{-pair} \text{ then } \text{partial-vector-derivative } (\lambda a. F\ a \cdot i) \ j\ a \text{ else } 0)$

$0) =$
 $(\lambda x. \text{if } x \in \text{Dy-pair} \text{ then if } x \in \text{Dy-pair} \text{ then partial-vector-derivative}$
 $(\lambda a. F a \cdot i) j x \text{ else } 0 \text{ else } 0)$
by auto
have $\text{arg2}: \text{Dy-pair} = \{(x, y). (\forall i \in \text{Basis}. a \cdot i \leq x \cdot i \wedge x \cdot i \leq b \cdot i) \wedge$
 $(\forall i \in \text{Basis}. g2 x \cdot i \leq y \cdot i \wedge y \cdot i \leq g1 x \cdot i)\}$
using Dy-def by auto
have $\text{arg3}: \bigwedge x. x \in \text{Dy-pair} \implies (\lambda x. \text{if } x \in \text{Dy-pair} \text{ then partial-vector-derivative}$
 $(\lambda a. F a \cdot i) j x \text{ else } 0) x$
 $= (\lambda x. \text{partial-vector-derivative } (\lambda a. F a \cdot i) j x) x$
by auto
have $\text{arg4}: \bigwedge x. x \in (\text{cbox } a \ b) \implies$
 $(\lambda x. \text{integral } (\text{cbox } (g2 \ x) \ (g1 \ x)) (\lambda y. \text{if } (x, y) \in \text{Dy-pair}$
 $\text{then partial-vector-derivative } (\lambda a. F a \cdot i) j (x, y) \text{ else } 0)) x =$
 $(\lambda x. \text{integral } (\text{cbox } (g2 \ x) \ (g1 \ x)) (\lambda y.$
 $\text{partial-vector-derivative } (\lambda a. F a \cdot i) j (x, y))) x$
apply (simp add: Dy-def)
by (smt Henstock-Kurzweil-Integration.integral-cong atLeastAtMost-iff)
show ?thesis
using gauge-integral-Fubini-curve-bounded-region-x
 $[\text{OF } f\text{-lesbegue-integrable } x\text{-axis-gauge-integrable } x\text{-axis-integral-measurable}$
 $\text{arg arg2}]$
 $\text{Henstock-Kurzweil-Integration.integral-cong}[\text{OF arg3, of Dy-pair } (\lambda x. x)]$
 $\text{Henstock-Kurzweil-Integration.integral-cong}[\text{OF arg4, of cbox } a \ b (\lambda x. x)]$
by auto
qed
have $5: (\text{integral Dy-pair } (\lambda a. (?F\text{-}b' \ a)))$
 $= \text{integral } (\text{cbox } a \ b) (\lambda x. F(x, g1(x)) \cdot i - F(x, g2(x)) \cdot i)$
using $4 \text{ Henstock-Kurzweil-Integration.integral-cong partial-deriv-one-d-integrable}$
 integrable-def
by (smt integral-unique)
show $(\text{line-integral } F \ \{i\} \ \text{gamma1}) + (\text{line-integral } F \ \{i\} \ \text{gamma2}) -$
 $(\text{line-integral } F \ \{i\} \ \text{gamma3}) - (\text{line-integral } F \ \{i\} \ \text{gamma4}) -$
 $= (\text{integral Dy-pair } (\lambda a. - (?F\text{-}b' \ a)))$
using $0(1) \ 1(1) \ 2(1) \ 3(1) \ 5 \ 6 \ 7$
by (simp add: Henstock-Kurzweil-Integration.integral-diff)
qed

theorem Greens-thm-type-II:

fixes $F:: ((\text{real} * \text{real}) \Rightarrow (\text{real} * \text{real}))$ **and**
 $\text{gamma4 } \text{gamma3 } \text{gamma2 } \text{gamma1} :: (\text{real} \Rightarrow (\text{real} * \text{real}))$ **and**
 $a:: \text{real}$ **and** $b:: \text{real}$ **and**
 $g1:: (\text{real} \Rightarrow \text{real})$ **and** $g2:: (\text{real} \Rightarrow \text{real})$
assumes $\text{Dx-def}: \text{Dx-pair} = \{(x::\text{real}, y) . y \in \text{cbox } a \ b \wedge x \in \text{cbox } (g2 \ y) \ (g1 \ y)\}$
and
 $\text{gamma4-def}: \text{gamma4} = (\lambda x. (g2(a + (b - a) * x), a + (b - a) * x))$ **and**
 $\text{gamma4-smooth}: \text{gamma4} \text{ piecewise-C1-differentiable-on } \{0..1\}$ **and**
 $\text{gamma3-def}: \text{gamma3} = (\lambda x. (g2(b) + x *_{\text{R}} (g1(b) - g2(b)), b))$ **and**
 $\text{gamma2-def}: \text{gamma2} = (\lambda x. (g1(a + (b - a) * x), a + (b - a) * x))$ **and**

gamma2-smooth: gamma2 piecewise-C1-differentiable-on {0..1} and
*gamma1-def: gamma1 = (λx. (g2(a) + x *_R (g1(a) - g2(a)), a)) and*
F-j-analytically-valid: analytically-valid Dx-pair (λp. F(p) · j) i and
g2-leq-g1: ∀ x ∈ cbox a b. (g2 x) ≤ (g1 x) and
a-lt-b: a < b
shows $-(\text{line-integral } F \{j\} \text{ gamma4}) -$
 $(\text{line-integral } F \{j\} \text{ gamma3}) +$
 $(\text{line-integral } F \{j\} \text{ gamma2}) +$
 $(\text{line-integral } F \{j\} \text{ gamma1})$
 $= (\text{integral } Dx\text{-pair } (\lambda a. (\text{partial-vector-derivative } (\lambda a. (F a) \cdot j) \ i$
a)))
line-integral-exists F {j} gamma4
line-integral-exists F {j} gamma3
line-integral-exists F {j} gamma2
line-integral-exists F {j} gamma1
proof –
let $?F\text{-}a' = \text{partial-vector-derivative } (\lambda a. (F a) \cdot j) \ i$
have *F-first-is-continuous: continuous-on Dx-pair (λa. F(a) · j)*
using *F-j-analytically-valid*
by *(auto simp add: analytically-valid-def)*
let $?f = (\lambda x. \text{if } x \in (Dx\text{-pair}) \text{ then } (\text{partial-vector-derivative } (\lambda a. (F a) \cdot j) \ i)$
x else 0)
have *f-lesbegue-integrable: integrable lborel ?f*
using *F-j-analytically-valid*
by *(auto simp add: analytically-valid-def indic-ident)*
have *partially-vec-diff: ∀ a ∈ Dx-pair. partially-vector-differentiable (λa. (F a) ·*
j) i a
using *F-j-analytically-valid*
by *(auto simp add: analytically-valid-def indicator-def)*
have $\bigwedge x \ y. ?F\text{-}a' (x,y) * \text{indicator } (Dx\text{-pair}) (x,y) = (\lambda x. \text{if } x \in (Dx\text{-pair}) \text{ then}$
 $?F\text{-}a' \ x \text{ else } 0) (x,y)$
using *indic-ident[of ?F-a'] by auto*
then have *y-axis-integral-measurable: (λx. integral UNIV (λy. ?f(y, x))) ∈*
borel-measurable lborel
using *analytically-valid-y[OF F-j-analytically-valid]*
by *(auto simp add: indicator-def)*
have *F-partially-differentiable: ∀ a ∈ Dx-pair. has-partial-vector-derivative (λx. (F*
x) · j) i (?F-a' a) a
using *partial-vector-derivative-works partially-vec-diff by fastforce*
have *g1-g2-continuous: continuous-on (cbox a b) g1 continuous-on (cbox a b) g2*
proof –
have *shift-scale-cont: continuous-on {a..b} (λx. (x - a)*(1/(b-a)))*
by *(intro continuous-intros)*
have *shift-scale-inv: (λx. a + (b - a) * x) ∘ (λx. (x - a)*(1/(b-a))) = id*
using *a-lt-b by (auto simp add: o-def)*
have *img-shift-scale:*
 $(\lambda x. (x - a)*(1/(b-a))) \{a..b\} = \{0..1\}$
apply *(auto simp: divide-simps image-iff)*
apply *(rule-tac x=x * (b - a) + a in bexI)*

```

    using a-lt-b by (auto simp: algebra-simps mult-le-cancel-left affine-ineq)
    have continuous-on {0..1} (λx. g2(a + (b - a) * x))
    using continuous-on-inner[OF piecewise-C1-differentiable-on-imp-continuous-on[OF
gamma4-smooth], of (λx. i), OF continuous-on-const]
    by (simp add: gamma4-def i-is-x-axis)
    then have continuous-on {a..b} ((λx. g2(a + (b - a) * x)) ∘ (λx. (x -
a)*(1/(b-a))))
    using img-shift-scale continuous-on-compose shift-scale-cont by force
    then show continuous-on (cbox a b) g2
    using a-lt-b by (simp add: shift-scale-inv)
    have continuous-on {0..1} (λx. g1(a + (b - a) * x))
    using continuous-on-inner[OF piecewise-C1-differentiable-on-imp-continuous-on[OF
gamma2-smooth], of (λx. i), OF continuous-on-const]
    by (simp add: gamma2-def i-is-x-axis)
    then have continuous-on {a..b} ((λx. g1(a + (b - a) * x)) ∘ (λx. (x -
a)*(1/(b-a))))
    using img-shift-scale continuous-on-compose shift-scale-cont by force
    then show continuous-on (cbox a b) g1
    using a-lt-b by (simp add: shift-scale-inv)
qed
have g2-scale-i-contin: continuous-on (cbox a b) (λx. (g2 x, 0))
  by (intro continuous-intros g1-g2-continuous)
let ?Dg2 = {p. ∃x. x ∈ cbox a b ∧ p = (g2(x), x)}
have line-is-pair-img: ?Dg2 = (λx. (g2(x), x)) ‘ (cbox a b)
  using image-def by auto
have g2-path-continuous: continuous-on (cbox a b) (λx. (g2(x), x))
  by (intro continuous-intros g1-g2-continuous)
have field-cont-on-gamma4-image: continuous-on ?Dg2 (λa. F(a) · j)
  by (rule continuous-on-subset [OF F-first-is-continuous]) (auto simp: Dx-def
g2-leq-g1)
have gamma4-is-compos-of-scal-and-g2: gamma4 = (λx. (g2(x), x)) ∘ (λx. a +
(b - a) * x)
  using gamma4-def by auto
have add-scale-img:
  (λx. a + (b - a) * x) ‘ {0 .. 1} = {a .. b} using add-scale-img and a-lt-b by
auto
then have Dg2-is-gamma4-pathimg: path-image gamma4 = ?Dg2
  using line-is-pair-img and gamma4-is-compos-of-scal-and-g2 image-comp path-image-def
  by (metis (no-types, lifting) cbox-interval)
have Base-vecs: i ∈ Basis j ∈ Basis i ≠ j
  using real-pair-basis and i-is-x-axis and j-is-y-axis by auto
have gamma4-as-euclid-space-fun: gamma4 = (λx. (a + (b - a) * x) *R j + (g2
(a + (b - a) * x), 0))
  using j-is-y-axis gamma4-def
  by auto
have 0: (line-integral F {j} gamma4) = integral (cbox a b) (λx. F(g2(x), x) · j)
  line-integral-exists F {j} gamma4
  using line-integral-on-pair-path-strong [OF norm-Basis[OF Base-vecs(2)] -
gamma4-as-euclid-space-fun]

```

```

    gamma4-def gamma4-smooth g2-scale-i-contin a-lt-b add-scale-img
    Dg2-is-gamma4-pathimg and field-cont-on-gamma4-image
  by (auto simp: pathstart-def pathfinish-def j-is-y-axis)
  then show line-integral-exists F {j} gamma4 by metis
  have gamma3-y-const:  $\forall x. \text{gamma3 } x \cdot j = b$ 
  by (simp add: gamma3-def j-is-y-axis)
  have 1: (line-integral F {j} gamma3) = 0 (line-integral-exists F {j} gamma3)
  using line-integral-on-pair-straight-path[OF gamma3-y-const] straight-path-differentiable-y
  gamma3-def
  by (auto simp add: mult.commute)
  then show line-integral-exists F {j} gamma3 by auto
  have continuous-on (cbox a b) ( $\lambda x. F(g2(x), x) \cdot j$ )
  by (smt Collect-mono-iff continuous-on-compose2 continuous-on-eq field-cont-on-gamma4-image
  g2-path-continuous line-is-pair-img)
  then have 6: ( $\lambda x. F(g2(x), x) \cdot j$ ) integrable-on (cbox a b)
  using integrable-continuous by blast
  have g1-scale-i-contin: continuous-on (cbox a b) ( $\lambda x. (g1 x, 0)$ )
  by (intro continuous-intros g1-g2-continuous)
  let ?Dg1 = {p.  $\exists x. x \in \text{cbox } a \ b \wedge p = (g1(x), x)$ }
  have line-is-pair-img: ?Dg1 = ( $\lambda x. (g1(x), x)$ ) ‘ (cbox a b)
  using image-def by auto
  have g1-path-continuous: continuous-on (cbox a b) ( $\lambda x. (g1(x), x)$ )
  by (intro continuous-intros g1-g2-continuous)
  have field-cont-on-gamma2-image: continuous-on ?Dg1 ( $\lambda a. F(a) \cdot j$ )
  by (rule continuous-on-subset [OF F-first-is-continuous]) (auto simp: Dx-def
  g2-leq-g1)
  have gamma2 = ( $\lambda x. (g1(x), x)$ )  $\circ$  ( $\lambda x. a + (b - a) * x$ )
  using gamma2-def by auto
  then have Dg1-is-gamma2-pathimg: path-image gamma2 = ?Dg1
  using line-is-pair-img image-comp path-image-def add-scale-img
  by (metis (no-types, lifting) cbox-interval)
  have Base-vecs:  $i \in \text{Basis } j \in \text{Basis } i \neq j$ 
  using real-pair-basis and i-is-x-axis and j-is-y-axis by auto
  have gamma2-as-euclid-space-fun: gamma2 = ( $\lambda x. (a + (b - a) * x) *_R j + (g1$ 
  ( $a + (b - a) * x$ ), 0))
  using j-is-y-axis gamma2-def by auto
  have 2: (line-integral F {j} gamma2) = integral (cbox a b) ( $\lambda x. F(g1(x), x) \cdot j$ )
  (line-integral-exists F {j} gamma2)
  using line-integral-on-pair-path-strong [OF norm-Basis[OF Base-vecs(2)] -
  gamma2-as-euclid-space-fun]
  gamma2-def and gamma2-smooth and g1-scale-i-contin and a-lt-b and
  add-scale-img
  Dg1-is-gamma2-pathimg and field-cont-on-gamma2-image
  by (auto simp: pathstart-def pathfinish-def j-is-y-axis)
  then show line-integral-exists F {j} gamma2 by metis
  have gamma1-y-const:  $\forall x. \text{gamma1 } x \cdot j = a$ 
  using gamma1-def
  by (auto simp add: real-inner-class.inner-add-left
  euclidean-space-class.inner-not-same-Basis j-is-y-axis)

```



```

have 3: (line-integral F {j} gamma1) = 0 (line-integral-exists F {j} gamma1)
using line-integral-on-pair-straight-path[OF gamma1-y-const] straight-path-differentiable-y
gamma1-def
by (auto simp add: mult.commute)
then show line-integral-exists F {j} gamma1 by auto
have continuous-on (cbox a b) (λx. F(g1(x), x) · j)
by (smt Collect-mono-iff continuous-on-compose2 continuous-on-eq field-cont-on-gamma2-image
g1-path-continuous line-is-pair-img)
then have 7: (λx. F(g1(x), x) · j) integrable-on (cbox a b)
using integrable-continuous [of a b (λx. F(g1(x), x) · j)]
by auto
have partial-deriv-one-d-integrable:
((λy. ?F-a'(y, xc)) has-integral F(g1(xc), xc) · j - F(g2(xc), xc) · j) (cbox (g2
xc) (g1 xc))
if xc ∈ cbox a b for xc::real
proof -
have {(y, xc'). y ∈ cbox (g2 xc) (g1 xc) ∧ xc' = xc} ⊆ Dx-pair
using that by (auto simp add: Dx-def)
then show ?thesis
using that and Base-vecs and F-partially-differentiable
and Dx-def [symmetric] and g2-leq-g1
and fundamental-theorem-of-calculus-partial-vector
[of g2 xc g1 xc i j xc *R j Dx-pair F ?F-a']
by (auto simp add: Groups.ab-semigroup-add-class.add commute i-is-x-axis
j-is-y-axis)
qed
have ?f integrable-on UNIV
by (simp add: f-lesbeque-integrable integrable-on-lborel)
then have partial-deriv-integrable: ?F-a' integrable-on Dx-pair
using integrable-restrict-UNIV by auto
have 4: integral Dx-pair ?F-a' = integral (cbox a b) (λx. integral (cbox (g2 x)
(g1 x)) (λy. ?F-a'(y, x)))
proof -
have y-axis-gauge-integrable: (λy. ?f(y, x)) integrable-on UNIV for x
proof -
let ?F-a'-oneD = (λx. (λy. if y ∈ (cbox (g2 x) (g1 x)) then ?F-a'(y, x) else
0))
have ∀x. x ∉ cbox a b → (λy. ?f(y, x)) = (λy. 0)
by (auto simp add: Dx-def)
then have f-integrable-x-not-in-range:
∀x. x ∉ cbox a b → (λy. ?f(y, x)) integrable-on UNIV
by (simp add: integrable-0)
have ∀x ∈ cbox a b. ?F-a'-oneD x = (λy. ?f(y, x))
using g2-leq-g1 by (auto simp add: Dx-def)
moreover have ∀x ∈ cbox a b. ?F-a'-oneD x integrable-on UNIV
using has-integral-integrable integrable-restrict-UNIV partial-deriv-one-d-integrable
by blast
ultimately have ∀x. x ∈ cbox a b → (λy. ?f(y, x)) integrable-on UNIV
by auto

```

then show $(\lambda y. ?f (y, x))$ *integrable-on UNIV*
using *f-integrable-x-not-in-range* **by auto**
qed
have *arg*: $(\lambda a. \text{if } a \in \text{Dx-pair} \text{ then } \text{partial-vector-derivative } (\lambda a. F a \cdot j) i a \text{ else } 0) =$
 $(\lambda x. \text{if } x \in \text{Dx-pair} \text{ then } \text{if } x \in \text{Dx-pair} \text{ then } \text{partial-vector-derivative } (\lambda a. F a \cdot j) i x \text{ else } 0 \text{ else } 0)$
by auto
have *arg2*: $\text{Dx-pair} = \{(y, x). (\forall i \in \text{Basis}. a \cdot i \leq x \cdot i \wedge x \cdot i \leq b \cdot i) \wedge (\forall i \in \text{Basis}. g2 x \cdot i \leq y \cdot i \wedge y \cdot i \leq g1 x \cdot i)\}$
using *Dx-def* **by auto**
have *arg3*: $\bigwedge x. x \in \text{Dx-pair} \implies (\lambda x. \text{if } x \in \text{Dx-pair} \text{ then } \text{partial-vector-derivative } (\lambda a. F a \cdot j) i x \text{ else } 0) x$
 $= (\lambda x. \text{partial-vector-derivative } (\lambda a. F a \cdot j) i x) x$
by auto
have *arg4*: $\bigwedge x. x \in (\text{cbox } a \ b) \implies$
 $(\lambda x. \text{integral } (\text{cbox } (g2 \ x) \ (g1 \ x)) (\lambda y. \text{if } (y, x) \in \text{Dx-pair} \text{ then } \text{partial-vector-derivative } (\lambda a. F a \cdot j) i (y, x) \text{ else } 0)) x =$
 $(\lambda x. \text{integral } (\text{cbox } (g2 \ x) \ (g1 \ x)) (\lambda y. \text{partial-vector-derivative } (\lambda a. F a \cdot j) i (y, x))) x$
apply *(clarsimp simp: Dx-def)*
by *(smt Henstock-Kurzweil-Integration.integral-cong atLeastAtMost-iff)*
show *?thesis*
using *gauge-integral-Fubini-curve-bounded-region-y*
 $[OF \text{ f-lesbeque-integrable } y\text{-axis-gauge-integrable } y\text{-axis-integral-measurable}$
 $\text{arg } \text{arg2}]$
 $\text{Henstock-Kurzweil-Integration.integral-cong}[OF \text{ arg3, of } \text{Dx-pair } (\lambda x. x)]$
 $\text{Henstock-Kurzweil-Integration.integral-cong}[OF \text{ arg4, of } \text{cbox } a \ b (\lambda x. x)]$
by auto
qed
have $((\text{integral } \text{Dx-pair } (\lambda a. (?F\text{-}a' \ a)))$
 $= \text{integral } (\text{cbox } a \ b) (\lambda x. F(g1(x), x) \cdot j - F(g2(x), x) \cdot j))$
using *4 Henstock-Kurzweil-Integration.integral-cong partial-deriv-one-d-integrable*
integrable-def
by *(smt integral-unique)*
then have $\text{integral } \text{Dx-pair } (\lambda a. - (?F\text{-}a' \ a))$
 $= - \text{integral } (\text{cbox } a \ b) (\lambda x. F(g1(x), x) \cdot j - F(g2(x), x) \cdot j)$
using *partial-deriv-integrable and integral-neg* **by auto**
then have *5*: $\text{integral } \text{Dx-pair } (\lambda a. - (?F\text{-}a' \ a))$
 $= \text{integral } (\text{cbox } a \ b) (\lambda x. (F(g2(x), x) \cdot j - F(g1(x), x) \cdot j))$
using *6 7*
and *integral-neg* $[of - (\lambda x. F(g1 \ x, x) \cdot j - F(g2 \ x, x) \cdot j)]$ **by auto**
have $(\text{line-integral } F \ \{j\} \ \text{gamma4}) + (\text{line-integral } F \ \{j\} \ \text{gamma3}) -$
 $(\text{line-integral } F \ \{j\} \ \text{gamma2}) - (\text{line-integral } F \ \{j\} \ \text{gamma1})$
 $= (\text{integral } \text{Dx-pair } (\lambda a. - (?F\text{-}a' \ a)))$
using *0 1 2 3 5 6 7*
 $\text{Henstock-Kurzweil-Integration.integral-diff}[of (\lambda x. F(g2(x), x) \cdot j)$
 $\text{cbox } a \ b (\lambda x. F(g1(x), x) \cdot j)]$ **by** *(auto)*
then show $-(\text{line-integral } F \ \{j\} \ \text{gamma4}) -$

$$\begin{aligned}
& (\text{line-integral } F \{j\} \text{ gamma3}) + \\
& (\text{line-integral } F \{j\} \text{ gamma2}) + \\
& (\text{line-integral } F \{j\} \text{ gamma1}) \\
& = (\text{integral } Dx\text{-pair } (\lambda a. (?F\text{-}a' a))) \\
\text{by } & (\text{simp add: } \langle \text{integral } Dx\text{-pair } (\lambda a. - ?F\text{-}a' a) = - \text{integral } (cbox a b) (\lambda x. \\
& F(g1 x, x) \cdot j - F(g2 x, x) \cdot j) \rangle \langle \text{integral } Dx\text{-pair } ?F\text{-}a' = \text{integral } (cbox a b) (\lambda x. \\
& F(g1 x, x) \cdot j - F(g2 x, x) \cdot j) \rangle) \\
\text{qed}
\end{aligned}$$

end

locale *green-typeII-cube* = *R2* +
fixes *twoC F*
assumes
two-cube: typeII-twoCube twoC and
valid-two-cube: valid-two-cube twoC and
f-analytically-valid: analytically-valid (cubeImage twoC) ($\lambda x. (F x) \cdot j$) i
begin

lemma *GreenThm-typeII-twoCube*:

shows *integral (cubeImage twoC) ($\lambda a. \text{partial-vector-derivative } (\lambda x. (F x) \cdot j$) i a) = one-chain-line-integral F {j} (boundary twoC)*
 $\forall (k, \gamma) \in \text{boundary twoC}. \text{line-integral-exists } F \{j\} \gamma$

proof –

let *?bottom-edge* = ($\lambda x. \text{twoC}(x, 0)$)
let *?right-edge* = ($\lambda y. \text{twoC}(1, y)$)
let *?top-edge* = ($\lambda x. \text{twoC}(x, 1)$)
let *?left-edge* = ($\lambda y. \text{twoC}(0, y)$)

have *line-integral-around-boundary*:

one-chain-line-integral F {j} (boundary twoC) =
line-integral F {j} ?bottom-edge + line-integral F {j} ?right-edge
– line-integral F {j} ?top-edge – line-integral F {j} ?left-edge

proof (*simp add: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def*)

have *finite1: finite {(- 1::int, $\lambda y. \text{twoC}(0, y)$), (1, $\lambda y. \text{twoC}(1, y)$), (- 1, $\lambda x. \text{twoC}(x, 1)$)}* **by** *auto*

then have *sum-step1: ($\sum (k, g) \in \{(- 1::int, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y)), (1, \lambda x. \text{twoC}(x, 0)), (- 1, \lambda x. \text{twoC}(x, 1))\}. k * \text{line-integral } F \{j\} g$)*
 $=$

*line-integral F {j} ($\lambda x. \text{twoC}(x, 0)$) + ($\sum (k, g) \in \{(- 1::int, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y)), (- 1, \lambda x. \text{twoC}(x, 1))\}. k * \text{line-integral } F \{j\} g$)*

using *sum.insert-remove [OF finite1]*

using *valid-two-cube*

apply (*simp only: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def valid-two-cube-def*)

by (*auto simp: card-insert-if split: if-split-asm*)

have *three-distinct-edges: card {(- 1::int, $\lambda y. \text{twoC}(0, y)$), (1, $\lambda y. \text{twoC}(1, y)$), (- 1, $\lambda x. \text{twoC}(x, 1)$)}* = 3

```

using valid-two-cube
apply (simp add: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def valid-two-cube-def)
by (auto simp: card-insert-if split: if-split-asm)
have finite2: finite  $\{(- 1::int, \lambda y. twoC (0, y)), (1, \lambda y. twoC (1, y))\}$  by auto
then have sum-step2:  $(\sum (k, g) \in \{(- (1::int), \lambda y. twoC (0, y)), (1, \lambda y. twoC (1, y)), (-1, \lambda x. twoC (x, 1))\}. k * line-integral F \{j\} g) =$ 
 $(- line-integral F \{j\} (\lambda x. twoC (x, 1))) + (\sum (k, g) \in \{(- (1::int), \lambda y. twoC (0, y)), (1, \lambda y. twoC (1, y))\}. k * line-integral F \{j\} g)$ 
using sum.insert-remove [OF finite2] three-distinct-edges
by (auto simp: card-insert-if split: if-split-asm)
have two-distinct-edges: card  $\{(- 1::int, \lambda y. twoC (0, y)), (1, \lambda y. twoC (1, y))\} = 2$ 
using three-distinct-edges
by (simp add: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def)
have finite3: finite  $\{(- 1::int, \lambda y. twoC (0, y))\}$  by auto
then have sum-step3:  $(\sum (k, g) \in \{(- (1::int), \lambda y. twoC (0, y)), (1, \lambda y. twoC (1, y))\}. k * line-integral F \{j\} g) =$ 
 $(line-integral F \{j\} (\lambda y. twoC (1, y))) + (\sum (k, g) \in \{(- (1::real), \lambda y. twoC (0, y))\}. k * line-integral F \{j\} g)$ 
using sum.insert-remove [OF finite2] three-distinct-edges
by (auto simp: card-insert-if split: if-split-asm)
show  $(\sum x \in \{(- 1::int, \lambda y. twoC (0, y)), (1, \lambda y. twoC (1, y)), (1, \lambda x. twoC (x, 0)), (- 1, \lambda x. twoC (x, 1))\}. case x of (k, g) \Rightarrow k * line-integral F \{j\} g) =$ 
 $line-integral F \{j\} (\lambda x. twoC (x, 0)) + line-integral F \{j\} (\lambda y. twoC (1, y)) - line-integral F \{j\} (\lambda x. twoC (x, 1)) - line-integral F \{j\} (\lambda y. twoC (0, y))$ 
using sum-step1 sum-step2 sum-step3 by auto
qed
obtain a b g1 g2 where
twoCisTypeII:  $a < b$ 
 $(\forall x \in cbox a b. g2 x \leq g1 x)$ 
cubeImage twoC =  $\{(y, x). x \in cbox a b \wedge y \in cbox (g2 x) (g1 x)\}$ 
twoC =  $(\lambda(y, x). ((1 - y) * g2 ((1 - x) * a + x * b) + y * g1 ((1 - x) * a + x * b), (1 - x) * a + x * b))$ 
g1 piecewise-C1-differentiable-on  $\{a .. b\}$ 
g2 piecewise-C1-differentiable-on  $\{a .. b\}$ 
using two-cube and typeII-twoCubeImg[oftwoC]
by auto
have left-edge-explicit:  $?left-edge = (\lambda x. (g2 (a + (b - a) * x), a + (b - a) * x))$ 
by (simp add: twoCisTypeII(4) algebra-simps)
have left-edge-smooth:  $?left-edge$  piecewise-C1-differentiable-on  $\{0..1\}$ 
proof -
have  $\forall x. (\lambda x. (a + (b - a) * x)) - ' \{x\} = \{(x - a)/(b - a)\}$ 
using twoCisTypeII(1) by (auto simp add: Set.vimage-def)
then have finite-vimg:  $\bigwedge x. finite(\{0..1\} \cap (\lambda x. (a + (b - a) * x)) - ' \{x\})$  by
auto

```

have *scale-shif-smth*: $(\lambda x. (a + (b - a) * x))$ *C1-differentiable-on* $\{0..1\}$ **using**
scale-shift-smooth **by** *auto*
then have *scale-shif-pw-smth*: $(\lambda x. (a + (b - a) * x))$ *piecewise-C1-differentiable-on*
 $\{0..1\}$ **using** *C1-differentiable-imp-piecewise* **by** *blast*
have *g2-smooth*: $g2$ *piecewise-C1-differentiable-on* $(\lambda x. a + (b - a) * x)$ ‘ $\{0..1\}$ ’
using *add-scale-img*[*OF twoCisTypeII(1)*] *twoCisTypeII(6)* **by** *auto*
have $(\lambda x. g2 (a + (b - a) * x))$ *piecewise-C1-differentiable-on* $\{0..1\}$
using *piecewise-C1-differentiable-compose*[*OF scale-shif-pw-smth g2-smooth*
finite-vimg]
by (*auto simp add: o-def*)
then have $(\lambda x::real. (g2 (a + (b - a) * x), a + (b - a) * x))$ *piece-*
wise-C1-differentiable-on $\{0..1\}$
using *all-components-smooth-one-pw-smooth-is-pw-smooth*[**where** $f = (\lambda x::real.$
 $(g2 (a + (b - a) * x), a + (b - a) * x))$]
by (*fastforce simp add: real-pair-basis*)
then show *?thesis* **using** *left-edge-explicit* **by** *auto*
qed
have *top-edge-explicit*: $?top-edge = (\lambda x. (g2 b + x *_R (g1 b - g2 b), b))$
and *right-edge-explicit*: $?right-edge = (\lambda x. (g1 (a + (b - a) * x), a + (b - a)$
 $* x))$
by (*simp-all add: twoCisTypeII(4) algebra-simps*)
have *right-edge-smooth*: $?right-edge$ *piecewise-C1-differentiable-on* $\{0..1\}$
proof –
have $\forall x. (\lambda x. (a + (b - a) * x)) - ' \{x\} = \{(x - a)/(b - a)\}$
using *twoCisTypeII(1)* **by** (*auto simp add: Set.vimage-def*)
then have *finite-vimg*: $\bigwedge x. \text{finite}(\{0..1\} \cap (\lambda x. (a + (b - a) * x)) - ' \{x\})$ **by**
auto
then have *scale-shif-pw-smth*: $(\lambda x. (a + (b - a) * x))$ *piecewise-C1-differentiable-on*
 $\{0..1\}$
using *C1-differentiable-imp-piecewise* [*OF scale-shift-smooth*] **by** *auto*
have *g1-smooth*: $g1$ *piecewise-C1-differentiable-on* $(\lambda x. a + (b - a) * x)$ ‘ $\{0..1\}$ ’
using *add-scale-img*[*OF twoCisTypeII(1)*] *twoCisTypeII(5)* **by** *auto*
have $(\lambda x. g1 (a + (b - a) * x))$ *piecewise-C1-differentiable-on* $\{0..1\}$
using *piecewise-C1-differentiable-compose*[*OF scale-shif-pw-smth g1-smooth*
finite-vimg]
by (*auto simp add: o-def*)
then have $(\lambda x::real. (g1 (a + (b - a) * x), a + (b - a) * x))$ *piece-*
wise-C1-differentiable-on $\{0..1\}$
using *all-components-smooth-one-pw-smooth-is-pw-smooth*[**where** $f = (\lambda x::real.$
 $(g1 (a + (b - a) * x), a + (b - a) * x))$]
by (*fastforce simp add: real-pair-basis*)
then show *?thesis* **using** *right-edge-explicit* **by** *auto*
qed
have *bottom-edge-explicit*: $?bottom-edge = (\lambda x. (g2 a + x *_R (g1 a - g2 a), a))$
by (*simp add: twoCisTypeII(4) algebra-simps*)
show *integral (cubeImage twoC)* $(\lambda a. \text{partial-vector-derivative } (\lambda x. (F x) \cdot j) i$
 $a) = \text{one-chain-line-integral } F \{j\}$ (*boundary twoC*)
using *Greens-thm-type-II*[*OF twoCisTypeII(3)*] *left-edge-explicit* *left-edge-smooth*
top-edge-explicit *right-edge-explicit* *right-edge-smooth*

```

    bottom-edge-explicit f-analytically-valid
    twoCisTypeII(2) twoCisTypeII(1)]
    line-integral-around-boundary
  by auto
  have line-integral-exists F {j}  $\gamma$  if  $(k, \gamma) \in \text{boundary twoC}$  for  $k \ \gamma$ 
  proof -
    have line-integral-exists F {j} ( $\lambda y. \text{twoC } (0, y)$ )
      line-integral-exists F {j} ( $\lambda x. \text{twoC } (x, 1)$ )
      line-integral-exists F {j} ( $\lambda y. \text{twoC } (1, y)$ )
      line-integral-exists F {j} ( $\lambda x. \text{twoC } (x, 0)$ )
    using Greens-thm-type-II[OF twoCisTypeII(3) left-edge-explicit left-edge-smooth
      top-edge-explicit right-edge-explicit right-edge-smooth
      bottom-edge-explicit f-analytically-valid
      twoCisTypeII(2) twoCisTypeII(1)] line-integral-around-boundary
    by auto
  then show line-integral-exists F {j}  $\gamma$ 
    using that by (auto simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
  qed
  then show  $\forall (k, \gamma) \in \text{boundary twoC}. \text{line-integral-exists F } \{j\} \ \gamma$  by auto
  qed

```

```

lemma line-integral-exists-on-typeII-Cube-boundaries':
  assumes  $(k, \gamma) \in \text{boundary twoC}$ 
  shows line-integral-exists F {j}  $\gamma$ 
  using assms GreenThm-typeII-twoCube(2) by blast

```

end

```

locale green-typeII-chain = R2 +
  fixes F two-chain s
  assumes valid-typeII-div: valid-typeII-division s two-chain and
    F-anal-valid:  $\forall \text{twoC} \in \text{two-chain}. \text{analytically-valid } (\text{cubeImage twoC}) (\lambda x. (F x) \cdot j) \ i$ 
  begin

```

```

lemma two-chain-valid-valid-cubes:  $\forall \text{two-cube} \in \text{two-chain}. \text{valid-two-cube two-cube}$ 
using valid-typeII-div
  by (auto simp add: valid-two-chain-def)

```

```

lemma typeII-chain-line-integral-exists-boundary':
  shows  $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain}. \text{line-integral-exists F } \{j\} \ \gamma$ 
  proof -
    have integ-exis:  $\forall (k, \gamma) \in \text{two-chain-boundary two-chain}. \text{line-integral-exists F } \{j\} \ \gamma$ 
      using green-typeII-cube.line-integral-exists-on-typeII-Cube-boundaries'[of i j]
      F-anal-valid valid-typeII-div
      apply (auto simp add: two-chain-boundary-def)

```

using *R2-axioms green-typeII-cube-axioms-def green-typeII-cube-def two-chain-valid-valid-cubes*
by *blast*
then show *integ-exis-vert:*
 $\forall (k,\gamma) \in \text{two-chain-vertical-boundary two-chain. line-integral-exists } F \{j\} \gamma$
by (*simp add: two-chain-boundary-def two-chain-vertical-boundary-def boundary-def*)
qed

lemma *typeII-chain-line-integral-exists-boundary'':*
 $\forall (k,\gamma) \in \text{two-chain-horizontal-boundary two-chain. line-integral-exists } F \{j\} \gamma$
proof –
have *integ-exis:* $\forall (k,\gamma) \in \text{two-chain-boundary two-chain. line-integral-exists } F \{j\} \gamma$
using *green-typeII-cube.line-integral-exists-on-typeII-Cube-boundaries'[of i j] valid-typeII-div*
apply (*simp add: two-chain-boundary-def boundary-def*)
using *F-anal-valid R2-axioms green-typeII-cube-axioms-def green-typeII-cube-def two-chain-valid-valid-cubes* **by** *fastforce*
then show *integ-exis-vert:*
 $\forall (k,\gamma) \in \text{two-chain-horizontal-boundary two-chain. line-integral-exists } F \{j\} \gamma$
by (*simp add: two-chain-boundary-def two-chain-horizontal-boundary-def boundary-def*)
qed

lemma *typeII-cube-line-integral-exists-boundary:*
 $\forall (k,\gamma) \in \text{two-chain-boundary two-chain. line-integral-exists } F \{j\} \gamma$
using *valid-typeII-div typeII-chain-line-integral-exists-boundary' typeII-chain-line-integral-exists-boundary''*
apply (*auto simp add: two-chain-boundary-def two-chain-horizontal-boundary-def two-chain-vertical-boundary-def*)
using *boundary-def* **by** *auto*

lemma *type-II-chain-horiz-bound-valid:*
 $\forall (k,\gamma) \in \text{two-chain-horizontal-boundary two-chain. valid-path } \gamma$
using *valid-typeII-div typeII-edges-are-valid-paths*
by (*force simp add: two-chain-boundary-def two-chain-horizontal-boundary-def boundary-def*)

lemma *type-II-chain-vert-bound-valid:*
 $\forall (k,\gamma) \in \text{two-chain-vertical-boundary two-chain. valid-path } \gamma$
using *typeII-edges-are-valid-paths valid-typeII-div*
by (*force simp add: two-chain-boundary-def two-chain-vertical-boundary-def boundary-def*)

lemma *members-of-only-horiz-div-line-integrable':*
assumes *only-horizontal-division one-chain two-chain*
 $(k::\text{int}, \gamma) \in \text{one-chain}$
 $(k::\text{int}, \gamma) \in \text{one-chain}$
finite two-chain
 $\forall \text{two-cube} \in \text{two-chain. valid-two-cube two-cube}$

shows *line-integral-exists* $F \{j\} \gamma$
proof –
have *integ-exis*: $\forall (k, \gamma) \in \text{two-chain-boundary two-chain. line-integral-exists } F \{j\} \gamma$
using *typeII-cube-line-integral-exists-boundary* **by** *blast*
have *integ-exis-vert*:
 $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain. line-integral-exists } F \{j\} \gamma$
using *typeII-chain-line-integral-exists-boundary'* *assms* **by** *auto*
have *integ-exis-horiz*:
 $(\bigwedge k \gamma. (\exists (k', \gamma') \in \text{two-chain-horizontal-boundary two-chain. } \exists a \in \{0..1\}. \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma) \implies \text{line-integral-exists } F \{j\} \gamma)$
using *integ-exis type-II-chain-horiz-bound-valid*
using *line-integral-exists-subpath*[of $F \{j\}$]
by (*fastforce simp add: two-chain-boundary-def two-chain-horizontal-boundary-def two-chain-vertical-boundary-def boundary-def*)
obtain $\mathcal{H} \ \mathcal{V}$ **where** *hv-props: finite* \mathcal{H}
 $(\forall (k, \gamma) \in \mathcal{H}. (\exists (k', \gamma') \in \text{two-chain-horizontal-boundary two-chain. } (\exists a \in \{0 .. 1\}. \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma)))$
 $\text{one-chain} = \mathcal{H} \cup \mathcal{V}$
 $((\text{common-sudiv-exists } (\text{two-chain-vertical-boundary two-chain}) \ \mathcal{V}) \vee \text{common-reparam-exists } \mathcal{V} \ (\text{two-chain-vertical-boundary two-chain}))$
finite \mathcal{V}
boundary-chain \mathcal{V}
 $\forall (k, \gamma) \in \mathcal{V}. \text{valid-path } \gamma$
using *assms(1) unfolding only-horizontal-division-def* **by** *blast*
have *finite-j*: *finite* $\{j\}$ **by** *auto*
show *line-integral-exists* $F \{j\} \gamma$
proof (*cases common-sudiv-exists (two-chain-vertical-boundary two-chain) \mathcal{V}*)
case *True*
show *?thesis*
using *gen-common-subdivision-imp-eq-line-integral(2)*[OF *True two-chain-vertical-boundary-is-boundary-ch*
hv-props(6) integ-exis-vert finite-two-chain-vertical-boundary[OF *assms(4)*] *hv-props(5)*
finite-j]
integ-exis-horiz[of γ] *assms(3) case-prod-conv hv-props(2) hv-props(3)*
by *fastforce*
next
case *False*
have *i*: $\{j\} \subseteq \text{Basis}$ **using** *j-is-y-axis real-pair-basis* **by** *auto*
have *ii*: $\forall (k2, \gamma2) \in \text{two-chain-vertical-boundary two-chain. } \forall b \in \{j\}. \text{continuous-on } (\text{path-image } \gamma2) \ (\lambda x. F \ x \cdot b)$
using *F-anal-valid field-cont-on-typeII-region-cont-on-edges valid-typeII-div*
by (*fastforce simp add: analytically-valid-def two-chain-vertical-boundary-def boundary-def path-image-def*)
show *line-integral-exists* $F \{j\} \gamma$
using *common-reparam-exists-imp-eq-line-integral(2)*[OF *finite-j hv-props(5)*
finite-two-chain-vertical-boundary[OF *assms(4)*] *hv-props(6) two-chain-vertical-boundary-is-boundary-ch*
ii - hv-props(7) type-II-chain-vert-bound-valid]
integ-exis-horiz[of γ] *assms(3) hv-props False*

by fastforce
qed
qed

lemma *GreenThm-typeII-twoChain:*

shows *two-chain-integral two-chain (partial-vector-derivative ($\lambda a. (F a) \cdot j$) i)*
= *one-chain-line-integral F {j} (two-chain-boundary two-chain)*

proof (*simp add: two-chain-boundary-def one-chain-line-integral-def two-chain-integral-def*)

let $?F\text{-}a' = \text{partial-vector-derivative } (\lambda a. (F a) \cdot j) \ i$

have $(\sum (k,g) \in \text{boundary twoCube}. k * \text{line-integral } F \{j\} \ g) = \text{integral } (\text{cubeImage twoCube}) \ (\lambda a. ?F\text{-}a' \ a)$

if *twoCube* \in *two-chain* for *twoCube*

using *green-typeII-cube.GreenThm-typeII-twoCube(1) valid-typeII-div F-anal-valid one-chain-line-integral-def valid-two-chain-def*

by (*simp add: R2-axioms green-typeII-cube-axioms-def green-typeII-cube-def that*)

then have *double-sum-eq-sum:*

$(\sum \text{twoCube} \in (\text{two-chain}). (\sum (k,g) \in \text{boundary twoCube}. k * \text{line-integral } F \{j\} \ g))$
= $(\sum \text{twoCube} \in (\text{two-chain}). \text{integral } (\text{cubeImage twoCube}) \ (\lambda a. ?F\text{-}a' \ a))$

using *Finite-Cartesian-Product.sum-cong-aux* by *auto*

have *pairwise-disjoint-boundaries:* $\forall x \in (\text{boundary 'two-chain}). (\forall y \in (\text{boundary 'two-chain}). (x \neq y \longrightarrow (x \cap y = \{\})))$

using *valid-typeII-div* by (*fastforce simp add: image-def valid-two-chain-def pairwise-def*)

have *finite-boundaries:* $\forall B \in (\text{boundary 'two-chain}). \text{finite } B$

using *valid-typeII-div image-iff* by (*fastforce simp add: valid-two-cube-def valid-two-chain-def*)

have *boundary-inj:* *inj-on boundary two-chain*

using *valid-typeII-div* by (*force simp add: valid-two-cube-def valid-two-chain-def pairwise-def inj-on-def*)

have $(\sum x \in (\bigcup x \in \text{two-chain}. \text{boundary } x). \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{j\} \ g) =$
 $(\sum \text{twoCube} \in (\text{two-chain}). (\sum (k,g) \in \text{boundary twoCube}. k * \text{line-integral } F \{j\} \ g))$

using *sum.reindex[OF boundary-inj, of $\lambda x. (\sum (k,g) \in x. k * \text{line-integral } F \{j\} \ g)$]*

using *sum.Union-disjoint[OF finite-boundaries*

pairwise-disjoint-boundaries,

*of $\lambda x. \text{case } x \text{ of } (k, g) \Rightarrow (k::\text{int}) * \text{line-integral } F \{j\} \ g$]*

by *auto*

then show $(\sum C \in \text{two-chain}. \text{integral } (\text{cubeImage } C) \ (\lambda a. ?F\text{-}a' \ a)) = (\sum x \in (\bigcup x \in \text{two-chain}. \text{boundary } x). \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{j\} \ g)$

using *double-sum-eq-sum* by *auto*

qed

lemma *GreenThm-typeII-divisible:*

assumes

gen-division: gen-division s (cubeImage ‘ two-chain)
shows *integral s (partial-vector-derivative (λx. (F x) · j) i) = one-chain-line-integral F {j} (two-chain-boundary two-chain)*
proof –
let $?F\text{-}a' = (\text{partial-vector-derivative } (\lambda x. (F x) \cdot j) i)$
have *integral s (λx. ?F-a' x) = two-chain-integral two-chain (λa. ?F-a' a)*
proof (*simp add: two-chain-integral-def*)
have *partial-deriv-integrable:*
(?F-a' has-integral (integral (cubeImage twoCube) ?F-a')) (cubeImage twoCube)
if *twoCube ∈ two-chain for twoCube*
by (*simp add: analytically-valid-imp-part-deriv-integrable-on F-anal-valid-integrable-integral that*)
then have *partial-deriv-integrable:*
 $\bigwedge \text{twoCubeImg. twoCubeImg} \in \text{cubeImage ' two-chain} \implies ((\lambda x. ?F\text{-}a' x)$
has-integral (integral (twoCubeImg) (λx. ?F-a' x))) (twoCubeImg)
using *integrable-neg by force*
have *finite-images: finite (cubeImage ‘ two-chain)*
using *gen-division gen-division-def by auto*
have *negligible-images: pairwise (λS S'. negligible (S ∩ S')) (cubeImage ‘ two-chain)*
using *gen-division by (auto simp add: gen-division-def pairwise-def)*
have *inj-on cubeImage two-chain using valid-typeII-div valid-two-chain-def by auto*
then have $(\sum \text{twoCubeImg} \in \text{cubeImage ' two-chain. integral twoCubeImg } (\lambda x. ?F\text{-}a' x))$
 $= (\sum C \in \text{two-chain. integral (cubeImage C) } (\lambda a. ?F\text{-}a' a))$
using *sum.reindex by auto*
then show *integral s (λx. ?F-a' x) = (∑ C ∈ two-chain. integral (cubeImage C) (λa. ?F-a' a))*
using *has-integral-Union[OF finite-images partial-deriv-integrable negligible-images] gen-division*
by (*auto simp add: gen-division-def*)
qed
then show *?thesis*
using *GreenThm-typeII-twoChain F-anal-valid*
by *auto*
qed

lemma *GreenThm-typeII-divisible-region-boundary-gen:*
assumes *only-horizontal-division: only-horizontal-division γ two-chain*
shows *integral s (partial-vector-derivative (λx. (F x) · j) i) = one-chain-line-integral F {j} γ*
proof –
let $?F\text{-}a' = (\text{partial-vector-derivative } (\lambda x. (F x) \cdot j) i)$

have *horiz-line-integral-zero:*
one-chain-line-integral F {j} (two-chain-horizontal-boundary two-chain) = 0
proof (*simp add: one-chain-line-integral-def*)
have *line-integral F {j} (snd oneCube) = 0*

```

    if oneCube ∈ two-chain-horizontal-boundary(two-chain) for oneCube
  proof –
    from that obtain x y1 y2 k
      where horiz-edge-def: oneCube = (k, (λt::real. ((1 - t) * (y2) + t * y1,
x::real)))
      using valid-typeII-div
      by (auto simp add: typeII-twoCube-def two-chain-horizontal-boundary-def
horizontal-boundary-def)
      let ?horiz-edge = (snd oneCube)
      have horiz-edge-y-const: ∀ t. (?horiz-edge t) · j = x
        by (auto simp add: horiz-edge-def real-inner-class.inner-add-left
euclidean-space-class.inner-not-same-Basis j-is-y-axis)
      have horiz-edge-is-straight-path: ?horiz-edge = (λt. (y2 + t * (y1 - y2), x))
        by (auto simp: horiz-edge-def algebra-simps)
      have ∀ x. ?horiz-edge differentiable at x
        using horiz-edge-is-straight-path straight-path-differentiable-y
        by (metis mult-commute-abs)
      then show line-integral F {j} (snd oneCube) = 0
        using line-integral-on-pair-straight-path(1) j-is-y-axis real-pair-basis horiz-edge-y-const
        by blast
    qed
    then show (∑ x∈two-chain-horizontal-boundary two-chain. case x of (k, g) ⇒
k * line-integral F {j} g) = 0
      by (force intro: sum.neutral)
    qed

  have boundary-is-finite: finite (two-chain-boundary two-chain)
    unfolding two-chain-boundary-def
  proof (rule finite-UN-I)
    show finite two-chain
      using valid-typeII-div finite-image-iff gen-division-def valid-two-chain-def by
auto
    show ∧a. a ∈ two-chain ⇒ finite (boundary a)
      by (simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
    qed
  have boundary-is-vert-hor:
    two-chain-boundary two-chain =
      (two-chain-vertical-boundary two-chain) ∪
      (two-chain-horizontal-boundary two-chain)
  by (auto simp add: two-chain-boundary-def two-chain-vertical-boundary-def two-chain-horizontal-boundary-def
boundary-def)
  then have hor-vert-finite:
    finite (two-chain-vertical-boundary two-chain)
    finite (two-chain-horizontal-boundary two-chain)
  using boundary-is-finite by auto
  have horiz-verti-disjoint:
    (two-chain-vertical-boundary two-chain) ∩
    (two-chain-horizontal-boundary two-chain) = {}
  proof (simp add: two-chain-vertical-boundary-def two-chain-horizontal-boundary-def

```

horizontal-boundary-def
vertical-boundary-def
show $(\bigcup_{x \in \text{two-chain}} \{(-1, \lambda y. x(0, y)), (1::\text{int}, \lambda y. x(1::\text{real}, y))\}) \cap$
 $(\bigcup_{x \in \text{two-chain}} \{(1, \lambda z. x(z, 0)), (-1, \lambda z. x(z, 1))\}) = \{\}$
proof –
have $\{(-1, \lambda y. \text{twoCube}(0, y)), (1::\text{int}, \lambda y. \text{twoCube}(1, y))\} \cap$
 $\{(1, \lambda z. \text{twoCube2}(z, 0)), (-1, \lambda z. \text{twoCube2}(z, 1))\} = \{\}$
if $\text{twoCube} \in \text{two-chain}$ $\text{twoCube2} \in \text{two-chain}$ **for** twoCube twoCube2
proof (*cases twoCube = twoCube2*)
case *True*
have *card-4*: $\text{card} \{(-1, \lambda y. \text{twoCube2}(0::\text{real}, y)), (1::\text{int}, \lambda y. \text{twoCube2}(1, y)),$
 $(1, \lambda x. \text{twoCube2}(x, 0)), (-1, \lambda x. \text{twoCube2}(x, 1))\} = 4$
using *valid-typeII-div valid-two-chain-def* *that*(2)
by (*auto simp add: boundary-def vertical-boundary-def horizontal-boundary-def*
valid-two-cube-def)
show *?thesis*
using *card-4* **by** (*auto simp: True card-insert-if split: if-split-asm*)
next
case *False*
show *?thesis*
using *valid-typeII-div valid-two-chain-def*
by (*simp add: boundary-def vertical-boundary-def horizontal-boundary-def*
pairwise-def <twoCube ≠ twoCube2> that)
qed
then have $\bigcup ((\lambda \text{twoCube}. \{(-1, \lambda y. \text{twoCube}(0::\text{real}, y)), (1::\text{real}, \lambda y.$
 $\text{twoCube}(1::\text{real}, y))\}) \text{ ‘two-chain}$
 $\cap \bigcup ((\lambda \text{twoCube}. \{(1::\text{int}, \lambda z. \text{twoCube}(z,$
 $0::\text{real}), (-1, \lambda z. \text{twoCube}(z, 1::\text{real}))\}) \text{ ‘two-chain}$
 $= \{\}$
by (*fastforce simp add: Union-disjoint*)
then show *?thesis* **by** *force*
qed
qed
have *one-chain-line-integral* $F \{j\}$ (*two-chain-boundary two-chain*)
 $= \text{one-chain-line-integral } F \{j\} (\text{two-chain-vertical-boundary two-chain}) +$
 $\text{one-chain-line-integral } F \{j\} (\text{two-chain-horizontal-boundary two-chain})$
using *boundary-is-vert-hor horiz-verti-disjoint*
by (*simp add: hor-vert-finite sum.union-disjoint one-chain-line-integral-def*)
then have *y-axis-line-integral-is-only-vertical*:
 $\text{one-chain-line-integral } F \{j\} (\text{two-chain-boundary two-chain})$
 $= \text{one-chain-line-integral } F \{j\} (\text{two-chain-vertical-boundary}$
 $\text{two-chain})$
using *horiz-line-integral-zero*
by *auto*
obtain $\mathcal{H} \mathcal{V}$ **where** *hv-props: finite* \mathcal{H}
 $(\forall (k, \gamma) \in \mathcal{H}. (\exists (k', \gamma') \in \text{two-chain-horizontal-boundary two-chain}.$
 $(\exists a \in \{0 .. 1\}.$
 $\exists b \in \{0..1\}.$

$a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma))$

$\gamma = \mathcal{H} \cup \mathcal{V}$
(common-sudiv-exists (two-chain-vertical-boundary two-chain) \mathcal{V}
 \vee common-reparam-exists \mathcal{V} (two-chain-vertical-boundary two-chain))
finite \mathcal{V}
boundary-chain \mathcal{V}
 $\forall (k, \gamma) \in \mathcal{V}$. *valid-path γ*
using *only-horizontal-division*
by(*fastforce simp add: only-horizontal-division-def*)
have *finite $\{j\}$ by auto*
then have *eq-integrals: one-chain-line-integral $F \{j\} \mathcal{V} =$ one-chain-line-integral*
 $F \{j\}$ (two-chain-vertical-boundary two-chain)
proof(*cases common-sudiv-exists (two-chain-vertical-boundary two-chain) \mathcal{V}*)
case *True then show ?thesis*
using *gen-common-subdivision-imp-eq-line-integral(1)[OF True two-chain-vertical-boundary-is-boundary-ch*
hv-props(6) - hor-vert-finite(1) hv-props(5)]
typeII-chain-line-integral-exists-boundary'
by force
next
case *False*
have *integ-exis-vert:*
 $\forall (k, \gamma) \in$ *two-chain-vertical-boundary two-chain. line-integral-exists $F \{j\} \gamma$*
using *typeII-chain-line-integral-exists-boundary' assms*
by (*fastforce simp add: valid-two-chain-def*)
have *integ-exis: $\forall (k, \gamma) \in$ two-chain-boundary two-chain. line-integral-exists F*
 $\{j\} \gamma$
using *typeII-cube-line-integral-exists-boundary by blast*
have *valid-paths: $\forall (k, \gamma) \in$ two-chain-horizontal-boundary two-chain. valid-path*
 γ
using *typeII-edges-are-valid-paths valid-typeII-div*
by (*fastforce simp add: two-chain-boundary-def two-chain-horizontal-boundary-def*
boundary-def)
have *integ-exis-horiz:*
 $(\bigwedge k \ \gamma. (\exists (k', \gamma') \in$ *two-chain-horizontal-boundary two-chain. $\exists a \in \{0..1\}. \exists b \in \{0..1\}.$*
 $a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma) \implies$
line-integral-exists $F \{j\} \gamma$)
using *integ-exis valid-paths line-integral-exists-subpath[of $F \{j\}$]*
by (*fastforce simp add: two-chain-boundary-def two-chain-horizontal-boundary-def*
two-chain-vertical-boundary-def boundary-def)
have *finite-j: finite $\{j\}$ by auto*
have *i: $\{j\} \subseteq$ Basis using j-is-y-axis real-pair-basis by auto*
have *ii: $\forall (k2, \gamma2) \in$ two-chain-vertical-boundary two-chain. $\forall b \in \{j\}$. continu-*
ous-on (path-image $\gamma2$) $(\lambda x. F \ x \cdot b)$
using *valid-typeII-div field-cont-on-typeII-region-cont-on-edges F-anal-valid*
by (*fastforce simp add: analytically-valid-def two-chain-vertical-boundary-def*
boundary-def path-image-def)
show *one-chain-line-integral $F \{j\} \mathcal{V} =$ one-chain-line-integral $F \{j\}$ (two-chain-vertical-boundary*
two-chain)
using *hv-props(4) False common-reparam-exists-imp-eq-line-integral(1)[OF fi-*

nite-j hv-props(5) hor-vert-finite(1) hv-props(6) two-chain-vertical-boundary-is-boundary-chain
ii

- *hv-props(7) type-II-chain-vert-bound-valid]*

by *fastforce*

qed

have *line-integral-on-path:*

one-chain-line-integral F {j} γ =

one-chain-line-integral F {j} (two-chain-vertical-boundary two-chain)

proof (*simp only: one-chain-line-integral-def*)

have *line-integral F {j} (snd oneCube) = 0* if *oneCube: oneCube $\in \mathcal{H}$ for oneCube*

proof -

obtain *k γ where k-gamma: (k, γ) = oneCube*

using *prod.collapse* by *blast*

then obtain *k' γ' a b where kp-gammap:*

(k'::int, γ') \in two-chain-horizontal-boundary two-chain

a \in {0 .. 1}

b \in {0..1}

subpath a b γ' = γ

using *hv-props oneCube*

by (*smt case-prodE split-conv*)

obtain *x y1 y2 where horiz-edge-def: (k', γ') = (k', ($\lambda t::real. ((1 - t) * (y2) + t * y1, x::real))$)*

using *valid-typeII-div kp-gammap*

by (*auto simp add: typeII-twoCube-def two-chain-horizontal-boundary-def horizontal-boundary-def*)

have *horiz-edge-y-const: $\forall t. \gamma (t) \cdot j = x$*

using *horiz-edge-def kp-gammap(4)*

by (*auto simp add: real-inner-class.inner-add-left*

euclidean-space-class.inner-not-same-Basis j-is-y-axis subpath-def)

have *horiz-edge-is-straight-path:*

*$\gamma = (\lambda t::real. ((1*y2 - a*y2) + a*y1 + ((b-a)*y1 - (b - a)*y2)*t, x::real)$*

proof -

fix *t::real*

have *(1 - (b - a)*t + a) * (y2) + ((b-a)*t + a) * y1 = (1 - (b - a)*t + a) * (y2) + ((b-a)*t + a) * y1*

by *auto*

then have *$\gamma = (\lambda t::real. ((1 - (b - a)*t - a) * (y2) + ((b-a)*t + a) * y1, x::real)$*

using *horiz-edge-def Product-Type.snd-conv Product-Type.fst-conv kp-gammap(4)*

by (*simp add: subpath-def diff-diff-eq[symmetric]*)

also have *... = ($\lambda t::real. ((1*y2 - (b - a)*y2*t - a*y2) + ((b-a)*y1*t + a*y1), x::real)$*

by(*auto simp add: ring-class.ring-distrib(2) Groups.diff-diff-eq left-diff-distrib*)

also have *... = ($\lambda t::real. ((1*y2 - a*y2) + a*y1 + ((b-a)*y1 - (b - a)*y2)*t, x::real)$*

by (*force simp add: left-diff-distrib*)

finally show $\gamma = (\lambda t::real. ((1*y2 - a*y2) + a*y1 + ((b-a)*y1 - (b - a)*y2))*t, x::real)$.

qed

show $line_integral\ F\ \{j\}\ (snd\ oneCube) = 0$

proof –

have $\forall x. \gamma$ differentiable at x

by (*simp add: horiz-edge-is-straight-path straight-path-differentiable-y*)

then have $line_integral\ F\ \{j\}\ \gamma = 0$

by (*simp add: horiz-edge-y-const line-integral-on-pair-straight-path(1)*)

then show *?thesis*

using *Product-Type.snd-conv k-gamma* **by** *auto*

qed

qed

then have $\forall x \in \mathcal{H}. (case\ x\ of\ (k, g) \Rightarrow (k::int) * line_integral\ F\ \{j\}\ g) = 0$

by *auto*

then show $(\sum x \in \gamma. case\ x\ of\ (k, g) \Rightarrow real_of_int\ k * line_integral\ F\ \{j\}\ g) =$

$(\sum x \in two_chain_vertical_boundary\ two_chain. case\ x\ of\ (k, g) \Rightarrow$

$real_of_int\ k * line_integral\ F\ \{j\}\ g)$

using *hv-props(1) hv-props(3) hv-props(5) sum-zero-set hor-vert-finite(1) eq-integrals*

by (*clarsimp simp add: one-chain-line-integral-def sum-zero-set*)

qed

then have $one_chain_line_integral\ F\ \{j\}\ \gamma =$

$one_chain_line_integral\ F\ \{j\}\ (two_chain_vertical_boundary$

$two_chain)$

using *line-integral-on-path* **by** *auto*

then have $one_chain_line_integral\ F\ \{j\}\ \gamma =$

$one_chain_line_integral\ F\ \{j\}\ (two_chain_boundary\ two_chain)$

using *y-axis-line-integral-is-only-vertical* **by** *auto*

then show *?thesis*

using *valid-typeII-div GreenThm-typeII-divisible* **by** *auto*

qed

lemma *GreenThm-typeII-divisible-region-boundary:*

assumes

two-cubes-trace-vertical-boundaries:

two-chain-vertical-boundary two-chain $\subseteq \gamma$ **and**

boundary-of-region-is-subset-of-partition-boundary:

$\gamma \subseteq two_chain_boundary\ two_chain$

shows $integral\ s\ (partial_vector_derivative\ (\lambda x. (F\ x) \cdot j)\ i) = one_chain_line_integral\ F\ \{j\}\ \gamma$

proof –

let $?F_a' = (partial_vector_derivative\ (\lambda x. (F\ x) \cdot j)\ i)$

have *horiz-line-integral-zero:*

$one_chain_line_integral\ F\ \{j\}\ (two_chain_horizontal_boundary\ two_chain) = 0$

proof (*simp add: one-chain-line-integral-def*)

have $line_integral\ F\ \{j\}\ (snd\ oneCube) = 0$

if *one: oneCube* $\in two_chain_horizontal_boundary(two_chain)$ **for** *oneCube*

```

proof –
  obtain  $x\ y1\ y2\ k$  where horiz-edge-def:  $oneCube = (k, (\lambda t::real. ((1 - t) * (y2) + t * y1, x::real)))$ 
  using valid-typeII-div one
  by (auto simp add: typeII-twoCube-def two-chain-horizontal-boundary-def horizontal-boundary-def)
  let  $?horiz-edge = (snd\ oneCube)$ 
  have horiz-edge-y-const:  $\forall t. (?horiz-edge\ t) \cdot j = x$ 
  using horiz-edge-def
  by (auto simp add: real-inner-class.inner-add-left euclidean-space-class.inner-not-same-Basis j-is-y-axis)
  have horiz-edge-is-straight-path:
     $?horiz-edge = (\lambda t. (y2 + t * (y1 - y2), x))$ 
  by (simp add: add-diff-eq diff-add-eq mult.commute right-diff-distrib horiz-edge-def)
  show  $line-integral\ F\ \{j\}\ (snd\ oneCube) = 0$ 
  by (metis horiz-edge-is-straight-path horiz-edge-y-const line-integral-on-pair-straight-path(1) mult.commute straight-path-differentiable-y)
  qed
  then show  $(\sum_{x \in two-chain-horizontal-boundary\ two-chain} case\ x\ of\ (k, g) \Rightarrow k * line-integral\ F\ \{j\}\ g) = 0$ 
  by (simp add: prod.case-eq-if)
  qed

  have boundary-is-finite:  $finite\ (two-chain-boundary\ two-chain)$ 
  unfolding two-chain-boundary-def
  proof (rule finite-UN-I)
  show  $finite\ two-chain$ 
  using valid-typeII-div finite-image-iff gen-division-def valid-two-chain-def by auto
  show  $\bigwedge a. a \in two-chain \Longrightarrow finite\ (boundary\ a)$ 
  by (simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
  qed
  have boundary-is-vert-hor:
     $two-chain-boundary\ two-chain = (two-chain-vertical-boundary\ two-chain) \cup (two-chain-horizontal-boundary\ two-chain)$ 
  by (auto simp add: two-chain-boundary-def two-chain-vertical-boundary-def two-chain-horizontal-boundary-def boundary-def)
  then have hor-vert-finite:
     $finite\ (two-chain-vertical-boundary\ two-chain)$ 
     $finite\ (two-chain-horizontal-boundary\ two-chain)$ 
  using boundary-is-finite by auto
  have horiz-verti-disjoint:
     $(two-chain-vertical-boundary\ two-chain) \cap (two-chain-horizontal-boundary\ two-chain) = \{\}$ 
  proof (simp add: two-chain-vertical-boundary-def two-chain-horizontal-boundary-def horizontal-boundary-def vertical-boundary-def)
  show  $(\bigcup_{x \in two-chain} \{(-1, \lambda y. x\ (0, y)), (1::int, \lambda y. x\ (1::real, y))\}) \cap (\bigcup_{x \in two-chain} \{(1, \lambda z. x\ (z, 0)), (-1, \lambda z. x\ (z, 1))\}) = \{\}$ 

```



```

proof –
  have  $\{(-1, \lambda y. \text{twoCube } (0, y)), (1::\text{int}, \lambda y. \text{twoCube } (1, y))\} \cap$ 
     $\{(1, \lambda y. \text{twoCube2 } (y, 0)), (-1, \lambda y. \text{twoCube2 } (y, 1))\} = \{\}$ 
  if  $\text{twoCube} \in \text{two-chain } \text{twoCube2} \in \text{two-chain}$  for  $\text{twoCube } \text{twoCube2}$ 
  proof ( $\text{cases } \text{twoCube} = \text{twoCube2}$ )
    case True
      have  $\text{card } \{(-1, \lambda y. \text{twoCube2 } (0::\text{real}, y)), (1::\text{int}, \lambda y. \text{twoCube2 } (1, y)),$ 
         $(1, \lambda x. \text{twoCube2 } (x, 0)), (-1, \lambda x. \text{twoCube2 } (x, 1))\} = 4$ 
      using valid-typeII-div valid-two-chain-def that(2)
      by (auto simp add: boundary-def vertical-boundary-def horizontal-boundary-def
        valid-two-cube-def)
      then show ?thesis
      by (auto simp: True card-insert-if split: if-split-asm)
    next
      case False show ?thesis
      using valid-typeII-div valid-two-chain-def
      by (simp add: boundary-def vertical-boundary-def horizontal-boundary-def
        pairwise-def <twoCube ≠ twoCube2> that(1) that(2))
      qed
      then have  $\bigcup ((\lambda \text{twoCube}. \{(-1, \lambda y. \text{twoCube } (0::\text{real}, y)), (1::\text{real}, \lambda y.$ 
         $\text{twoCube } (1::\text{real}, y))\}) \text{ 'two-chain})$ 
         $\cap \bigcup ((\lambda \text{twoCube}. \{(1::\text{int}, \lambda z. \text{twoCube } (z,$ 
         $0::\text{real}), (-1, \lambda z. \text{twoCube } (z, 1::\text{real}))\}) \text{ 'two-chain})$ 
         $= \{\}$ 
      by (force simp add: Complete-Lattices.Union-disjoint)
      then show ?thesis by force
    qed
  qed
  have  $\text{one-chain-line-integral } F \{j\} (\text{two-chain-boundary } \text{two-chain})$ 
     $= \text{one-chain-line-integral } F \{j\} (\text{two-chain-vertical-boundary } \text{two-chain}) +$ 
     $\text{one-chain-line-integral } F \{j\} (\text{two-chain-horizontal-boundary } \text{two-chain})$ 
  using boundary-is-vert-hor horiz-verti-disjoint
  by (auto simp add: one-chain-line-integral-def hor-vert-finite sum.union-disjoint)
  then have y-axis-line-integral-is-only-vertical:
     $\text{one-chain-line-integral } F \{j\} (\text{two-chain-boundary } \text{two-chain})$ 
     $= \text{one-chain-line-integral } F \{j\} (\text{two-chain-vertical-boundary } \text{two-chain})$ 
  using horiz-line-integral-zero by auto

  have  $\exists \mathcal{H}. \mathcal{H} \subseteq (\text{two-chain-horizontal-boundary } \text{two-chain}) \wedge$ 
     $\gamma = \mathcal{H} \cup (\text{two-chain-vertical-boundary } \text{two-chain})$ 
  proof
    let  $? \mathcal{H} = \gamma - (\text{two-chain-vertical-boundary } \text{two-chain})$ 
    show  $? \mathcal{H} \subseteq \text{two-chain-horizontal-boundary } \text{two-chain} \wedge \gamma = ? \mathcal{H} \cup \text{two-chain-vertical-boundary}$ 
    two-chain
    using two-cubes-trace-vertical-boundaries
    boundary-of-region-is-subset-of-partition-boundary boundary-is-vert-hor
    by blast
  qed
  then obtain  $\mathcal{H}$  where

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    h-props:  $\mathcal{H} \subseteq (\text{two-chain-horizontal-boundary two-chain})$ 
     $\gamma = \mathcal{H} \cup (\text{two-chain-vertical-boundary two-chain})$ 
    by auto
  have h-vert-disj:  $\mathcal{H} \cap (\text{two-chain-vertical-boundary two-chain}) = \{\}$ 
    using horiz-verti-disjoint h-props(1) by auto
  have h-finite: finite  $\mathcal{H}$ 
    using hor-vert-finite h-props(1) Ffinite-Set.rev-finite-subset by force
  have line-integral-on-path:
    one-chain-line-integral  $F \{j\} \gamma =$ 
    one-chain-line-integral  $F \{j\} \mathcal{H} + \text{one-chain-line-integral } F \{j\} (\text{two-chain-vertical-boundary two-chain})$ 
    by (auto simp add: one-chain-line-integral-def h-props sum.union-disjoint[OF h-finite hor-vert-finite(1) h-vert-disj])

  have one-chain-line-integral  $F \{j\} \mathcal{H} = 0$ 
  proof (simp add: one-chain-line-integral-def)
    have line-integral  $F \{j\} (\text{snd oneCube}) = 0$ 
      if oneCube:  $\text{oneCube} \in \text{two-chain-horizontal-boundary}(\text{two-chain})$  for oneCube
    proof -
      obtain  $x \ y1 \ y2 \ k$  where horiz-edge-def:  $\text{oneCube} = (k, (\lambda t::\text{real}. ((1 - t) * (y2) + t * y1, x::\text{real})))$ 
        using valid-typeII-div oneCube
      by (auto simp add: typeII-twoCube-def two-chain-horizontal-boundary-def horizontal-boundary-def)
      let ?horiz-edge = (snd oneCube)
      have horiz-edge-y-const:  $\forall t. (?horiz-edge \ t) \cdot j = x$ 
        by (simp add: j-is-y-axis horiz-edge-def)
      have horiz-edge-is-straight-path:
        ?horiz-edge =  $(\lambda t. (y2 + t * (y1 - y2), x))$ 
      by (simp add: add-diff-eq diff-add-eq mult.commute right-diff-distrib horiz-edge-def)
      show line-integral  $F \{j\} (\text{snd oneCube}) = 0$ 
      by (simp add: horiz-edge-is-straight-path j-is-y-axis line-integral-on-pair-straight-path(1) mult.commute straight-path-differentiable-y)
    qed
  then have  $\forall \text{oneCube} \in \mathcal{H}. \text{line-integral } F \{j\} (\text{snd oneCube}) = 0$ 
    using h-props by auto
  then have  $\forall x \in \mathcal{H}. (\text{case } x \text{ of } (k, g) \Rightarrow (k::\text{int}) * \text{line-integral } F \{j\} g) = 0$ 
    by auto
  then show  $(\sum x \in \mathcal{H}. \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{j\} g) = 0$ 
    by simp
  qed
  then have one-chain-line-integral  $F \{j\} \gamma =$ 
    one-chain-line-integral  $F \{j\} (\text{two-chain-vertical-boundary two-chain})$ 
    using line-integral-on-path by auto
  then have one-chain-line-integral  $F \{j\} \gamma =$ 
    one-chain-line-integral  $F \{j\} (\text{two-chain-boundary two-chain})$ 
    using y-axis-line-integral-is-only-vertical by auto
  then show ?thesis

```

using *valid-typeII-div GreenThm-typeII-divisible* by *auto*
qed

end

locale *green-typeI-cube* = *R2* +

fixes *twoC F*

assumes

two-cube: *typeI-twoCube twoC* **and**

valid-two-cube: *valid-two-cube twoC* **and**

f-analytically-valid: *analytically-valid (cubeImage twoC) ($\lambda x. (F x) \cdot i$) j*

begin

lemma *GreenThm-typeI-twoCube*:

shows *integral (cubeImage twoC) ($\lambda a. - \text{partial-vector-derivative } (\lambda p. F p \cdot i) j$) a* = *one-chain-line-integral F {i} (boundary twoC)*

$\forall (k, \gamma) \in \text{boundary twoC}. \text{line-integral-exists } F \{i\} \gamma$

proof –

let *?bottom-edge* = ($\lambda x. \text{twoC}(x, 0)$)

let *?right-edge* = ($\lambda y. \text{twoC}(1, y)$)

let *?top-edge* = ($\lambda x. \text{twoC}(x, 1)$)

let *?left-edge* = ($\lambda y. \text{twoC}(0, y)$)

have *line-integral-around-boundary*:

one-chain-line-integral F {i} (boundary twoC) =

line-integral F {i} ?bottom-edge + *line-integral F {i} ?right-edge*
– *line-integral F {i} ?top-edge* – *line-integral F {i} ?left-edge*

proof (*simp add: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def*)

have *finite1*: *finite* $\{(- 1::\text{int}, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y)), (- 1, \lambda x. \text{twoC}(x, 1))\}$ **by** *auto*

have *sum-step1*: $(\sum (k, g) \in \{(- 1::\text{int}, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y)), (1, \lambda x. \text{twoC}(x, 0)), (- 1, \lambda x. \text{twoC}(x, 1))\}. k * \text{line-integral } F \{i\} g) =$
 $\text{line-integral } F \{i\} (\lambda x. \text{twoC}(x, 0)) + (\sum (k, g) \in \{(- 1::\text{int}), \lambda y. \text{twoC}(0, y), (1, \lambda y. \text{twoC}(1, y)), (- 1, \lambda x. \text{twoC}(x, 1))\}. k * \text{line-integral } F \{i\} g)$

using *sum.insert-remove [OF finite1] valid-two-cube*

by (*auto simp: horizontal-boundary-def vertical-boundary-def boundary-def valid-two-cube-def card-insert-if split: if-split-asm*)

have *three-distinct-edges*: *card* $\{(- 1::\text{int}, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y)), (- 1, \lambda x. \text{twoC}(x, 1))\} = 3$

using *valid-two-cube*

apply (*simp add: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def valid-two-cube-def*)

by (*auto simp: card-insert-if split: if-split-asm*)

have *finite2*: *finite* $\{(- 1::\text{int}, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y))\}$ **by** *auto*

have *sum-step2*: $(\sum (k, g) \in \{(- 1::\text{int}, \lambda y. \text{twoC}(0, y)), (1, \lambda y. \text{twoC}(1, y)), (- 1, \lambda x. \text{twoC}(x, 1))\}. k * \text{line-integral } F \{i\} g) =$
 $(- \text{line-integral } F \{i\} (\lambda x. \text{twoC}(x, 1))) + (\sum (k, g) \in \{(- 1::\text{int}), \lambda y. \text{twoC}(0, y), (1, \lambda y. \text{twoC}(1, y))\}. k * \text{line-integral } F \{i\} g)$

```

using sum.insert-remove [OF finite2] three-distinct-edges
by (auto simp: card-insert-if split: if-split-asm)
have two-distinct-edges:  $\text{card} \{(- 1::\text{int}, \lambda y. \text{twoC } (0, y)), (1, \lambda y. \text{twoC } (1, y))\} = 2$ 
using three-distinct-edges
by (simp add: one-chain-line-integral-def horizontal-boundary-def vertical-boundary-def boundary-def)
have finite3: finite  $\{(- 1::\text{int}, \lambda y. \text{twoC } (0, y))\}$  by auto
have sum-step3:  $(\sum (k, g) \in \{(- 1::\text{int}, \lambda y. \text{twoC } (0, y)), (1, \lambda y. \text{twoC } (1, y))\}. k * \text{line-integral } F \{i\} g) =$ 
 $(\text{line-integral } F \{i\} (\lambda y. \text{twoC } (1, y))) + (\sum (k, g) \in \{(- 1::\text{real}, \lambda y. \text{twoC } (0, y))\}. k * \text{line-integral } F \{i\} g)$ 
using sum.insert-remove [OF finite2] three-distinct-edges
by (auto simp: card-insert-if split: if-split-asm)
show  $(\sum x \in \{(- 1::\text{int}, \lambda y. \text{twoC } (0, y)), (1, \lambda y. \text{twoC } (1, y)), (1, \lambda x. \text{twoC } (x, 0)), (- 1, \lambda x. \text{twoC } (x, 1))\}. \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{i\} g) =$ 
 $\text{line-integral } F \{i\} (\lambda x. \text{twoC } (x, 0)) + \text{line-integral } F \{i\} (\lambda y. \text{twoC } (1, y)) - \text{line-integral } F \{i\} (\lambda x. \text{twoC } (x, 1)) - \text{line-integral } F \{i\} (\lambda y. \text{twoC } (0, y))$ 
using sum-step1 sum-step2 sum-step3 by auto
qed
obtain a b g1 g2 where
twoCisTypeI:  $a < b$ 
 $(\forall x \in \text{cbox } a \ b. g2 \ x \leq g1 \ x)$ 
cubeImage twoC =  $\{(x, y). x \in \text{cbox } a \ b \wedge y \in \text{cbox } (g2 \ x) \ (g1 \ x)\}$ 
twoC =  $(\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2 \ ((1 - x) * a + x * b) + y * g1 \ ((1 - x) * a + x * b)))$ 
g1 piecewise-C1-differentiable-on  $\{a .. b\}$ 
g2 piecewise-C1-differentiable-on  $\{a .. b\}$ 
 $(\lambda x. \text{twoC}(x, 0)) = (\lambda x. (a + (b - a) * x, g2 \ (a + (b - a) * x)))$ 
 $(\lambda y. \text{twoC}(1, y)) = (\lambda x. (b, g2 \ b + x *_R \ (g1 \ b - g2 \ b)))$ 
 $(\lambda x. \text{twoC}(x, 1)) = (\lambda x. (a + (b - a) * x, g1 \ (a + (b - a) * x)))$ 
 $(\lambda y. \text{twoC}(0, y)) = (\lambda x. (a, g2 \ a + x *_R \ (g1 \ a - g2 \ a)))$ 
using two-cube and typeI-cube-explicit-spec[oftwoC] by auto
have bottom-edge-smooth:  $(\lambda x. \text{twoC } (x, 0))$  piecewise-C1-differentiable-on  $\{0..1\}$ 
using typeI-twoCube-smooth-edges two-cube boundary-def vertical-boundary-def horizontal-boundary-def
by auto
have top-edge-smooth: ?top-edge piecewise-C1-differentiable-on  $\{0..1\}$ 
using typeI-twoCube-smooth-edges two-cube boundary-def vertical-boundary-def horizontal-boundary-def
by auto
show integral (cubeImage twoC)  $(\lambda a. - \text{partial-vector-derivative } (\lambda p. F \ p \cdot i) \ j \ a) = \text{one-chain-line-integral } F \{i\} \ (\text{boundary } \text{twoC})$ 
using Greens-thm-type-I[OF twoCisTypeI(3) twoCisTypeI(7) bottom-edge-smooth twoCisTypeI(8) twoCisTypeI(9) top-edge-smooth twoCisTypeI(10) f-analytically-valid twoCisTypeI(2) twoCisTypeI(1)]
line-integral-around-boundary
by auto

```

have *line-integral-exists* $F \{i\}$ ($\lambda y. \text{twoC } (0, y)$)
line-integral-exists $F \{i\}$ ($\lambda x. \text{twoC } (x, 1)$)
line-integral-exists $F \{i\}$ ($\lambda y. \text{twoC } (1, y)$)
line-integral-exists $F \{i\}$ ($\lambda x. \text{twoC } (x, 0)$)
using *Greens-thm-type-I*[*OF twoCisTypeI*(3) *twoCisTypeI*(7) *bottom-edge-smooth*
twoCisTypeI(8) *twoCisTypeI*(9) *top-edge-smooth*
twoCisTypeI(10) *f-analytically-valid twoCisTypeI*(2) *twoCisTypeI*(1)]
line-integral-around-boundary
by *auto*
then have *line-integral-exists* $F \{i\}$ γ **if** $(k, \gamma) \in \text{boundary twoC}$ **for** $k \gamma$
using *that by (auto simp add: boundary-def horizontal-boundary-def vertical-boundary-def)*
then show $\forall (k, \gamma) \in \text{boundary twoC}. \text{line-integral-exists } F \{i\} \gamma$ **by** *auto*
qed

lemma *line-integral-exists-on-typeI-Cube-boundaries'*:
assumes $(k, \gamma) \in \text{boundary twoC}$
shows *line-integral-exists* $F \{i\} \gamma$
using *assms two-cube valid-two-cube f-analytically-valid GreenThm-typeI-twoCube*(2)
by *blast*

end

locale *green-typeI-chain* = $R2 +$
fixes $F \text{ two-chain } s$
assumes *valid-typeI-div: valid-typeI-division s two-chain and*
F-anal-valid: $\forall \text{twoC} \in \text{two-chain}. \text{analytically-valid } (\text{cubeImage twoC}) (\lambda x. (F x) \cdot i) j$
begin

lemma *two-chain-valid-valid-cubes: $\forall \text{two-cube} \in \text{two-chain}. \text{valid-two-cube two-cube}$*
using *valid-typeI-div*
by *(auto simp add: valid-two-chain-def)*

lemma *typeI-cube-line-integral-exists-boundary'*:
assumes $\forall \text{two-cube} \in \text{two-chain}. \text{typeI-twoCube two-cube}$
assumes $\forall \text{twoC} \in \text{two-chain}. \text{analytically-valid } (\text{cubeImage twoC}) (\lambda x. (F x) \cdot i) j$
assumes $\forall \text{two-cube} \in \text{two-chain}. \text{valid-two-cube two-cube}$
shows $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain}. \text{line-integral-exists } F \{i\} \gamma$

γ

proof –

have *integ-exis: $\forall (k, \gamma) \in \text{two-chain-boundary two-chain}. \text{line-integral-exists } F \{i\} \gamma$*
using *green-typeI-cube.line-integral-exists-on-typeI-Cube-boundaries'*[*of i j*] *assms*
using *R2-axioms green-typeI-cube-axioms-def green-typeI-cube-def two-chain-boundary-def*
by *fastforce*
then show *integ-exis-vert:*
 $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain}. \text{line-integral-exists } F \{i\} \gamma$

by (*simp add: two-chain-boundary-def two-chain-vertical-boundary-def boundary-def*)
qed

lemma *typeI-cube-line-integral-exists-boundary''*:

$\forall (k, \gamma) \in \text{two-chain-horizontal-boundary two-chain. line-integral-exists } F \{i\} \gamma$

proof –

have *integ-exis*: $\forall (k, \gamma) \in \text{two-chain-boundary two-chain. line-integral-exists } F \{i\} \gamma$

using *green-typeI-cube.line-integral-exists-on-typeI-Cube-boundaries'[of i j] valid-typeI-div*
apply (*simp add: two-chain-boundary-def boundary-def*)

using *F-anal-valid R2-axioms green-typeI-cube-axioms-def green-typeI-cube-def two-chain-valid-valid-cubes* **by** *fastforce*

then show *integ-exis-horiz*:

$\forall (k, \gamma) \in \text{two-chain-horizontal-boundary two-chain. line-integral-exists } F \{i\} \gamma$

by (*simp add: two-chain-boundary-def two-chain-horizontal-boundary-def boundary-def*)

qed

lemma *typeI-cube-line-integral-exists-boundary*:

$\forall (k, \gamma) \in \text{two-chain-boundary two-chain. line-integral-exists } F \{i\} \gamma$

using *typeI-cube-line-integral-exists-boundary' typeI-cube-line-integral-exists-boundary''*

apply (*auto simp add: two-chain-boundary-def two-chain-horizontal-boundary-def two-chain-vertical-boundary-def*)

by (*meson R2-axioms green-typeI-chain.F-anal-valid green-typeI-chain-axioms green-typeI-cube.line-integral-exists-on-typeI-Cube-boundaries' green-typeI-cube-axioms-def green-typeI-cube-def two-chain-valid-valid-cubes valid-typeI-div*)

lemma *type-I-chain-horiz-bound-valid*:

$\forall (k, \gamma) \in \text{two-chain-horizontal-boundary two-chain. valid-path } \gamma$

using *typeI-edges-are-valid-paths valid-typeI-div*

by (*force simp add: two-chain-boundary-def two-chain-horizontal-boundary-def boundary-def*)

lemma *type-I-chain-vert-bound-valid*:

assumes $\forall \text{two-cube} \in \text{two-chain. typeI-twoCube two-cube}$

shows $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain. valid-path } \gamma$

using *typeI-edges-are-valid-paths assms(1)*

by (*force simp add: two-chain-boundary-def two-chain-vertical-boundary-def boundary-def*)

lemma *members-of-only-vertical-div-line-integrable'*:

assumes *only-vertical-division one-chain two-chain*

$(k::\text{int}, \gamma) \in \text{one-chain}$

$(k::\text{int}, \gamma) \in \text{one-chain}$

finite two-chain

shows *line-integral-exists* $F \{i\} \gamma$

proof –

have *integ-exis*: $\forall (k, \gamma) \in \text{two-chain-boundary two-chain. line-integral-exists } F$

```

{i}  $\gamma$ 
  using typeI-cube-line-integral-exists-boundary by blast
  have integ-exis-horiz:
     $\forall (k, \gamma) \in \text{two-chain-horizontal-boundary two-chain. line-integral-exists } F \{i\} \gamma$ 
    using typeI-cube-line-integral-exists-boundary'' assms by auto
  have valid-paths:  $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain. valid-path } \gamma$ 
    using type-I-chain-vert-bound-valid valid-typeI-div by linarith
  have integ-exis-vert:
     $(\bigwedge k \gamma. (\exists (k', \gamma') \in \text{two-chain-vertical-boundary two-chain. } \exists a \in \{0..1\}. \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma) \implies$ 
       $\text{line-integral-exists } F \{i\} \gamma)$ 
    using integ-exis valid-paths line-integral-exists-subpath[of  $F \{i\}$ ]
  by (fastforce simp add: two-chain-boundary-def two-chain-horizontal-boundary-def
    two-chain-vertical-boundary-def boundary-def)
  obtain  $\mathcal{H} \ \mathcal{V}$  where hv-props: finite  $\mathcal{V}$ 
     $(\forall (k, \gamma) \in \mathcal{V}. (\exists (k', \gamma') \in \text{two-chain-vertical-boundary two-chain.}$ 
       $(\exists a \in \{0..1\}. \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma)))$ 
    one-chain =  $\mathcal{H} \cup \mathcal{V}$ 
    (common-sudiv-exists (two-chain-horizontal-boundary two-chain)  $\mathcal{H}$ )
     $\vee$  common-reparam-exists  $\mathcal{H}$  (two-chain-horizontal-boundary two-chain)
    finite  $\mathcal{H}$ 
    boundary-chain  $\mathcal{H}$ 
     $\forall (k, \gamma) \in \mathcal{H}. \text{valid-path } \gamma$ 
  using assms(1) unfolding only-vertical-division-def by blast
  have finite-i: finite  $\{i\}$  by auto
  show line-integral-exists  $F \{i\} \gamma$ 
  proof (cases common-sudiv-exists (two-chain-horizontal-boundary two-chain)  $\mathcal{H}$ )
    case True
      show ?thesis
      using gen-common-subdivision-imp-eq-line-integral(2)[OF True two-chain-horizontal-boundary-is-boundary
        hv-props(6) integ-exis-horiz finite-two-chain-horizontal-boundary[OF assms(4)] hv-props(5)
        finite-i]
        integ-exis-vert[of  $\gamma$ ] assms(3) case-prod-conv hv-props(2) hv-props(3)
      by fastforce
    next
      case False
      have i:  $\{i\} \subseteq \text{Basis}$  using i-is-x-axis real-pair-basis by auto
      have ii:  $\forall (k2, \gamma2) \in \text{two-chain-horizontal-boundary two-chain. } \forall b \in \{i\}. \text{continuous-on}$ 
        (path-image  $\gamma2$ ) ( $\lambda x. F \ x \cdot b$ )
      using assms field-cont-on-typeI-region-cont-on-edges F-anal-valid valid-typeI-div
      by (fastforce simp add: analytically-valid-def two-chain-horizontal-boundary-def
        boundary-def path-image-def)
      show line-integral-exists  $F \{i\} \gamma$ 
      using common-reparam-exists-imp-eq-line-integral(2)[OF finite-i hv-props(5)
        finite-two-chain-horizontal-boundary[OF assms(4)] hv-props(6) two-chain-horizontal-boundary-is-boundary-ch
        ii
        - hv-props(7) type-I-chain-horiz-bound-valid]
      integ-exis-vert[of  $\gamma$ ] False
      assms(3) hv-props(2-4) by fastforce

```

qed
qed

lemma *GreenThm-typeI-two-chain:*

two-chain-integral two-chain $(\lambda a. - \text{partial-vector-derivative } (\lambda x. (F x) \cdot i) j a)$
 $= \text{one-chain-line-integral } F \{i\} (\text{two-chain-boundary two-chain})$

proof (*simp add: two-chain-boundary-def one-chain-line-integral-def two-chain-integral-def*)

let $?F\text{-}b' = \text{partial-vector-derivative } (\lambda x. (F x) \cdot i) j$

have *all-two-cubes-have-four-distict-edges:* $\forall \text{twoCube} \in \text{two-chain}. \text{card } (\text{boundary twoCube}) = 4$

using *valid-typeI-div valid-two-chain-def valid-two-cube-def* **by** *auto*

have *no-shared-edges-have-similar-orientations:*

$\forall \text{twoCube1} \in \text{two-chain}. \forall \text{twoCube2} \in \text{two-chain}. \text{twoCube1} \neq \text{twoCube2} \longrightarrow$
 $\text{boundary twoCube1} \cap \text{boundary twoCube2} = \{\}$

using *valid-typeI-div valid-two-chain-def*

by (*auto simp add: pairwise-def*)

have $(\sum (k,g) \in \text{boundary twoCube}. k * \text{line-integral } F \{i\} g) = \text{integral } (\text{cubeImage twoCube}) (\lambda a. - ?F\text{-}b' a)$

if $\text{twoCube} \in \text{two-chain}$ **for** twoCube

proof –

have *analytically-val:* *analytically-valid* $(\text{cubeImage twoCube}) (\lambda x. F x \cdot i) j$

using *that F-anal-valid* **by** *auto*

show $(\sum (k, g) \in \text{boundary twoCube}. k * \text{line-integral } F \{i\} g) = \text{integral } (\text{cubeImage twoCube}) (\lambda a. - ?F\text{-}b' a)$

using *green-typeI-cube.GreenThm-typeI-twoCube*

apply (*simp add: one-chain-line-integral-def*)

by (*simp add: R2-axioms analytically-val green-typeI-cube-axioms-def green-typeI-cube-def that two-chain-valid-valid-cubes valid-typeI-div*)

qed

then have *double-sum-eq-sum:*

$(\sum \text{twoCube} \in (\text{two-chain}). (\sum (k,g) \in \text{boundary twoCube}. k * \text{line-integral } F \{i\} g))$
 $= (\sum \text{twoCube} \in (\text{two-chain}). \text{integral } (\text{cubeImage twoCube}) (\lambda a. - ?F\text{-}b' a))$

using *Finite-Cartesian-Product.sum-cong-aux* **by** *auto*

have *pairwise-disjoint-boundaries:* $\forall x \in (\text{boundary 'two-chain}). (\forall y \in (\text{boundary 'two-chain}). (x \neq y \longrightarrow (x \cap y = \{\})))$

using *no-shared-edges-have-similar-orientations*

by (*force simp add: image-def disjoint-iff-not-equal*)

have *finite-boundaries:* $\forall B \in (\text{boundary 'two-chain}). \text{finite } B$

using *all-two-cubes-have-four-distict-edges*

using *image-iff* **by** *fastforce*

have *boundary-inj:* *inj-on* $\text{boundary two-chain}$

using *all-two-cubes-have-four-distict-edges* **and** *no-shared-edges-have-similar-orientations*

by (*force simp add: inj-on-def*)

have $(\sum x \in (\bigcup (\text{boundary 'two-chain})). \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{i\} g)$

$= (\text{sum} \circ \text{sum}) (\lambda x. \text{case } x \text{ of } (k, g) \Rightarrow (k::\text{int}) * \text{line-integral } F \{i\} g)$
 $(\text{boundary 'two-chain})$

using *sum.Union-disjoint*[*OF finite-boundaries pairwise-disjoint-boundaries*]
by *simp*
also have ... = $(\sum \text{twoCube} \in (\text{two-chain}). (\sum (k,g) \in \text{boundary twoCube}. k * \text{line-integral } F \{i\} g))$
using *sum.reindex*[*OF boundary-inj*, of $\lambda x. (\sum (k,g) \in x. k * \text{line-integral } F \{i\} g)$]
by *auto*
finally show $(\sum C \in \text{two-chain}. - \text{integral } (\text{cubeImage } C) (\text{partial-vector-derivative } (\lambda x. F x \cdot i) j)) = (\sum x \in (\bigcup x \in \text{two-chain}. \text{boundary } x). \text{case } x \text{ of } (k, g) \Rightarrow \text{real-of-int } k * \text{line-integral } F \{i\} g)$
using *double-sum-eq-sum* **by** *auto*
qed

lemma *GreenThm-typeI-divisible*:

assumes *gen-division*: *gen-division* s (*cubeImage* ' *two-chain*)
shows *integral* s ($\lambda x. - \text{partial-vector-derivative } (\lambda a. F(a) \cdot i) j x$) = *one-chain-line-integral* $F \{i\}$ (*two-chain-boundary* *two-chain*)
proof –
let $?F\text{-}b' = \text{partial-vector-derivative } (\lambda a. F(a) \cdot i) j$
have *integral* s ($\lambda x. - ?F\text{-}b' x$) = *two-chain-integral* *two-chain* ($\lambda a. - ?F\text{-}b' a$)
proof (*simp add*: *two-chain-integral-def*)
have $(\sum f \in \text{two-chain}. \text{integral } (\text{cubeImage } f) (\text{partial-vector-derivative } (\lambda p. F p \cdot i) j)) = \text{integral } s (\text{partial-vector-derivative } (\lambda p. F p \cdot i) j)$
by (*metis analytically-valid-imp-part-deriv-integrable-on F-anal-valid gen-division two-chain-integral-def two-chain-integral-eq-integral-divisible valid-typeI-div*)
then show – *integral* s (*partial-vector-derivative* ($\lambda a. F a \cdot i$) j) = $(\sum C \in \text{two-chain}. - \text{integral } (\text{cubeImage } C) (\text{partial-vector-derivative } (\lambda a. F a \cdot i) j))$
by (*simp add*: *sum-negf*)
qed
then show *?thesis*
using *GreenThm-typeI-two-chain assms* **by** *auto*
qed

lemma *GreenThm-typeI-divisible-region-boundary*:

assumes
gen-division: *gen-division* s (*cubeImage* ' *two-chain*) **and**
two-cubes-trace-horizontal-boundaries:
two-chain-horizontal-boundary *two-chain* $\subseteq \gamma$ **and**
boundary-of-region-is-subset-of-partition-boundary:
 $\gamma \subseteq \text{two-chain-boundary } \text{two-chain}$
shows *integral* s ($\lambda x. - \text{partial-vector-derivative } (\lambda a. F(a) \cdot i) j x$) = *one-chain-line-integral* $F \{i\}$ γ
proof –
let $?F\text{-}b' = \text{partial-vector-derivative } (\lambda a. F(a) \cdot i)$
have *all-two-cubes-have-four-distict-edges*: $\forall \text{twoCube} \in \text{two-chain}. \text{card } (\text{boundary } \text{twoCube}) = 4$
using *valid-typeI-div valid-two-chain-def valid-two-cube-def* **by** *auto*
have *no-shared-edges-have-similar-orientations*:
 $\forall \text{twoCube1} \in \text{two-chain}. \forall \text{twoCube2} \in \text{two-chain}.$

```

    twoCube1 ≠ twoCube2 → boundary twoCube1 ∩ boundary twoCube2 = {}
  using valid-typeI-div valid-two-chain-def by (auto simp add: pairwise-def)

  have vert-line-integral-zero:
    one-chain-line-integral F {i} (two-chain-vertical-boundary two-chain) = 0
  proof (simp add: one-chain-line-integral-def)
    have line-integral F {i} (snd oneCube) = 0
    if oneCube: oneCube ∈ two-chain-vertical-boundary(two-chain) for oneCube
    proof -
      obtain x y1 y2 k where vert-edge-def: oneCube = (k, (λt::real. (x::real, (1
      - t) * (y2) + t * y1)))
      using valid-typeI-div oneCube
      by (auto simp add: typeI-twoCube-def two-chain-vertical-boundary-def verti-
      cal-boundary-def)
      let ?vert-edge = (snd oneCube)
      have vert-edge-x-const: ∀ t. (?vert-edge t) · i = x
      by (simp add: i-is-x-axis vert-edge-def)
      have vert-edge-is-straight-path: ?vert-edge = (λt. (x, y2 + t * (y1 - y2)))
      using vert-edge-def Product-Type.snd-conv
      by (auto simp add: mult.commute right-diff-distrib')
      show ?thesis
      by (simp add: i-is-x-axis line-integral-on-pair-straight-path(1) mult.commute
      straight-path-differentiable-x vert-edge-is-straight-path)
    qed
    then show (∑ x∈two-chain-vertical-boundary two-chain. case x of (k, g) ⇒ k
    * line-integral F {i} g) = 0
    using comm-monoid-add-class.sum.neutral by (simp add: prod.case-eq-if)
  qed

  have boundary-is-finite: finite (two-chain-boundary two-chain)
    unfolding two-chain-boundary-def
    by (metis all-two-cubes-have-four-distict-edges card.infinite finite-UN-I finite-imageD

    gen-division gen-division-def zero-neq-numeral valid-typeI-div valid-two-chain-def)
  have boundary-is-vert-hor: (two-chain-boundary two-chain) =
    (two-chain-vertical-boundary two-chain) ∪
    (two-chain-horizontal-boundary two-chain)
  by (auto simp add: two-chain-boundary-def two-chain-vertical-boundary-def two-chain-horizontal-boundary-def
  boundary-def)
  then have hor-vert-finite:
    finite (two-chain-vertical-boundary two-chain)
    finite (two-chain-horizontal-boundary two-chain)
  using boundary-is-finite by auto
  have horiz-verti-disjoint:
    (two-chain-vertical-boundary two-chain) ∩ (two-chain-horizontal-boundary two-chain)
  = {}
  proof (simp add: two-chain-vertical-boundary-def two-chain-horizontal-boundary-def
  horizontal-boundary-def
  vertical-boundary-def)

```

show $(\bigcup x \in \text{two-chain}. \{(-1, \lambda y. x(0, y)), (1::\text{int}, \lambda y. x(1::\text{real}, y))\}) \cap$
 $(\bigcup x \in \text{two-chain}. \{(1, \lambda z. x(z, 0)), (-1, \lambda z. x(z, 1))\}) = \{\}$

proof –

have $\{(-1, \lambda y. \text{twoCube}(0, y)), (1::\text{int}, \lambda y. \text{twoCube}(1, y))\} \cap$
 $\{(1, \lambda z. \text{twoCube2}(z, 0)), (-1, \lambda z. \text{twoCube2}(z, 1))\} = \{\}$

if $\text{twoCube} \in \text{two-chain}$ $\text{twoCube2} \in \text{two-chain}$ **for** twoCube twoCube2

proof (*cases twoCube = twoCube2*)

case *True*

have $\text{card} \{(-1::\text{int}, \lambda y. \text{twoCube2}(0::\text{real}, y)), (1::\text{int}, \lambda y. \text{twoCube2}(1,$
 $y)), (1, \lambda x. \text{twoCube2}(x, 0)), (-1, \lambda x. \text{twoCube2}(x, 1))\} = 4$

using *all-two-cubes-have-four-distict-edges that(2)*

by (*auto simp add: boundary-def vertical-boundary-def horizontal-boundary-def*)

then show *?thesis*

by (*auto simp: True card-insert-if split: if-split-asm*)

next

case *False*

then show *?thesis*

using *no-shared-edges-have-similar-orientations*

by (*simp add: that boundary-def vertical-boundary-def horizontal-boundary-def*)

qed

then have $\bigcup ((\lambda \text{twoCube}. \{(-1::\text{int}, \lambda y. \text{twoCube}(0, y)), (1, \lambda y. \text{twoCube}$
 $(1, y))\}) \text{ ‘two-chain})$
 $\cap \bigcup ((\lambda \text{twoCube}. \{(1, \lambda y. \text{twoCube}(y, 0)), (-1, \lambda z. \text{twoCube}(z, 1))\})$
 $\text{ ‘two-chain}) = \{\}$

using *Complete-Lattices.Union-disjoint by force*

then show *?thesis by force*

qed

qed

have $\text{one-chain-line-integral } F \{i\} (\text{two-chain-boundary two-chain})$
 $= \text{one-chain-line-integral } F \{i\} (\text{two-chain-vertical-boundary two-chain}) +$
 $\text{one-chain-line-integral } F \{i\} (\text{two-chain-horizontal-boundary two-chain})$

using *boundary-is-vert-hor horiz-verti-disjoint*

by (*auto simp add: one-chain-line-integral-def hor-vert-finite sum.union-disjoint*)

then have *x-axis-line-integral-is-only-horizontal:*
 $\text{one-chain-line-integral } F \{i\} (\text{two-chain-boundary two-chain})$
 $= \text{one-chain-line-integral } F \{i\} (\text{two-chain-horizontal-boundary two-chain})$

using *vert-line-integral-zero by auto*

have $\exists \mathcal{V}. \mathcal{V} \subseteq (\text{two-chain-vertical-boundary two-chain}) \wedge \gamma = \mathcal{V} \cup (\text{two-chain-horizontal-boundary}$
 $\text{two-chain})$

proof

let $?V = \gamma - (\text{two-chain-horizontal-boundary two-chain})$

show $?V \subseteq \text{two-chain-vertical-boundary two-chain} \wedge \gamma = ?V \cup \text{two-chain-horizontal-boundary}$
 two-chain

using *two-cubes-trace-horizontal-boundaries*
boundary-of-region-is-subset-of-partition-boundary boundary-is-vert-hor

by *blast*

qed

then obtain \mathcal{V} **where**

$v\text{-props}: \mathcal{V} \subseteq (\text{two-chain-vertical-boundary two-chain}) \gamma = \mathcal{V} \cup (\text{two-chain-horizontal-boundary two-chain})$
by auto
have $v\text{-horiz-disj}: \mathcal{V} \cap (\text{two-chain-horizontal-boundary two-chain}) = \{\}$
using $\text{horiz-verti-disjoint } v\text{-props}(1)$ **by auto**
have $v\text{-finite}: \text{finite } \mathcal{V}$
using $\text{hor-vert-finite } v\text{-props}(1)$ $\text{Finite-Set.rev-finite-subset}$ **by force**
have $\text{line-integral-on-path}: \text{one-chain-line-integral } F \{i\} \gamma =$
 $\text{one-chain-line-integral } F \{i\} \mathcal{V} + \text{one-chain-line-integral}$
 $F \{i\} (\text{two-chain-horizontal-boundary two-chain})$
by $(\text{auto simp add: one-chain-line-integral-def } v\text{-props sum.union-disjoint}[OF$
 $v\text{-finite hor-vert-finite}(2) v\text{-horiz-disj}])$

have $\text{one-chain-line-integral } F \{i\} \mathcal{V} = 0$
proof $(\text{simp add: one-chain-line-integral-def})$
have $\text{line-integral } F \{i\} (\text{snd oneCube}) = 0$
if $\text{oneCube}: \text{oneCube} \in \text{two-chain-vertical-boundary}(\text{two-chain})$ **for** oneCube
proof –
obtain $x \ y1 \ y2 \ k$ **where** $\text{vert-edge-def}: \text{oneCube} = (k, (\lambda t::\text{real}. (x::\text{real}, (1$
 $- t) * (y2) + t * y1)))$
using $\text{valid-typeI-div oneCube}$
by $(\text{auto simp add: typeI-twoCube-def two-chain-vertical-boundary-def verti-}$
 $\text{cal-boundary-def})$
let $?vert\text{-edge} = (\text{snd oneCube})$
have $\text{vert-edge-x-const}: \forall t. (?vert\text{-edge } t) \cdot i = x$
by $(\text{simp add: i-is-x-axis vert-edge-def})$
have $\text{vert-edge-is-straight-path}: ?vert\text{-edge} = (\lambda t. (x, y2 + t * (y1 - y2)))$
by $(\text{auto simp: vert-edge-def algebra-simps})$
have $\forall x. ?vert\text{-edge}$ *differentiable at* x
by $(\text{metis mult.commute vert-edge-is-straight-path straight-path-differentiable-x})$
then show $\text{line-integral } F \{i\} (\text{snd oneCube}) = 0$
using $\text{line-integral-on-pair-straight-path}(1) \text{vert-edge-x-const}$ **by blast**
qed
then have $\forall \text{oneCube} \in \mathcal{V}. \text{line-integral } F \{i\} (\text{snd oneCube}) = 0$
using $v\text{-props}$ **by auto**
then show $(\sum x \in \mathcal{V}. \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{i\} g) = 0$
using $\text{comm-monoid-add-class.sum.neutral}$ **by** $(\text{simp add: prod.case-eq-if})$
qed
then have $\text{one-chain-line-integral } F \{i\} \gamma =$
 $\text{one-chain-line-integral } F \{i\} (\text{two-chain-boundary two-chain})$
using $x\text{-axis-line-integral-is-only-horizontal}$ **by** $(\text{simp add: line-integral-on-path})$
then show $?thesis$
using $\text{assms and GreenThm-typeI-divisible}$ **by auto**
qed

lemma $\text{GreenThm-typeI-divisible-region-boundary-gen}: \text{assumes } \text{valid-typeI-div}: \text{valid-typeI-division } s \text{ two-chain}$ **and**
 $f\text{-analytically-valid}: \forall \text{twoC} \in \text{two-chain}. \text{analytically-valid } (\text{cubeImage twoC})$

$(\lambda a. F(a) \cdot i) j$ **and**
only-vertical-division:
only-vertical-division γ two-chain
shows *integral s* $(\lambda x. - \text{partial-vector-derivative } (\lambda a. F(a) \cdot i) j x) = \text{one-chain-line-integral } F \{i\} \gamma$
proof –
let $?F\text{-}b' = \text{partial-vector-derivative } (\lambda a. F(a) \cdot i)$
have *all-two-cubes-have-four-distict-edges:* $\forall \text{twoCube} \in \text{two-chain}. \text{card } (\text{boundary } \text{twoCube}) = 4$
using *valid-typeI-div valid-two-chain-def valid-two-cube-def*
by *auto*
have *no-shared-edges-have-similar-orientations:*
 $\forall \text{twoCube1} \in \text{two-chain}. \forall \text{twoCube2} \in \text{two-chain}.$
 $\text{twoCube1} \neq \text{twoCube2} \longrightarrow \text{boundary } \text{twoCube1} \cap \text{boundary } \text{twoCube2} = \{\}$
using *valid-typeI-div valid-two-chain-def* **by** *(auto simp add: pairwise-def)*

have *vert-line-integral-zero:*
 $\text{one-chain-line-integral } F \{i\} (\text{two-chain-vertical-boundary } \text{two-chain}) = 0$
proof *(simp add: one-chain-line-integral-def)*
have $\text{line-integral } F \{i\} (\text{snd } \text{oneCube}) = 0$
if $\text{oneCube}: \text{oneCube} \in \text{two-chain-vertical-boundary}(\text{two-chain})$ **for** oneCube
proof –
obtain $x \ y1 \ y2 \ k$ **where** *vert-edge-def:* $\text{oneCube} = (k, (\lambda t::\text{real}. (x::\text{real}, (1 - t) * (y2) + t * y1)))$
using *valid-typeI-div oneCube*
by *(auto simp add: typeI-twoCube-def two-chain-vertical-boundary-def vertical-boundary-def)*
let $?vert\text{-}edge = (\text{snd } \text{oneCube})$
have *vert-edge-x-const:* $\forall t. (?vert\text{-}edge \ t) \cdot i = x$
by *(simp add: i-is-x-axis vert-edge-def)*
have *vert-edge-is-straight-path:* $?vert\text{-}edge = (\lambda t. (x, y2 + t * (y1 - y2)))$
by *(auto simp: vert-edge-def algebra-simps)*
show *?thesis*
by *(simp add: i-is-x-axis line-integral-on-pair-straight-path(1) mult.commute straight-path-differentiable-x vert-edge-is-straight-path)*
qed
then show $(\sum x \in \text{two-chain-vertical-boundary } \text{two-chain}. \text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{i\} g) = 0$
using *comm-monoid-add-class.sum.neutral* **by** *(simp add: prod.case-eq-if)*
qed

have *boundary-is-finite:* $\text{finite } (\text{two-chain-boundary } \text{two-chain})$
unfolding *two-chain-boundary-def*
proof *(rule finite-UN-I)*
show *finite two-chain*
using *assms(1) finite-imageD gen-division-def valid-two-chain-def* **by** *auto*
show $\bigwedge a. a \in \text{two-chain} \Longrightarrow \text{finite } (\text{boundary } a)$
by *(simp add: boundary-def horizontal-boundary-def vertical-boundary-def)*
qed

```

have boundary-is-vert-hor: two-chain-boundary two-chain =
      (two-chain-vertical-boundary two-chain)  $\cup$  (two-chain-horizontal-boundary
two-chain)
  by (auto simp add: two-chain-boundary-def two-chain-vertical-boundary-def
two-chain-horizontal-boundary-def boundary-def)
  then have hor-vert-finite:
    finite (two-chain-vertical-boundary two-chain)
    finite (two-chain-horizontal-boundary two-chain)
  using boundary-is-finite by auto
  have horiz-verti-disjoint:
    (two-chain-vertical-boundary two-chain)  $\cap$  (two-chain-horizontal-boundary two-chain)
= {}
  proof (simp add: two-chain-vertical-boundary-def two-chain-horizontal-boundary-def
horizontal-boundary-def
vertical-boundary-def)
  show ( $\bigcup x \in \text{two-chain}. \{(-1, \lambda y. x(0, y)), (1::\text{int}, \lambda y. x(1::\text{real}, y))\}$ )
 $\cap$  ( $\bigcup x \in \text{two-chain}. \{(1, \lambda y. x(y, 0)), (-1, \lambda y. x(y, 1))\}$ ) = {}
  proof -
  have  $\{(-1, \lambda y. \text{twoCube}(0, y)), (1::\text{int}, \lambda y. \text{twoCube}(1, y))\} \cap$ 
 $\{(1, \lambda y. \text{twoCube2}(y, 0)), (-1, \lambda y. \text{twoCube2}(y, 1))\} = \{\}$ 
  if twoCube  $\in$  two-chain twoCube2  $\in$  two-chain for twoCube twoCube2
  proof (cases twoCube = twoCube2)
  case True
  have card  $\{(-1::\text{int}, \lambda y. \text{twoCube2}(0, y)), (1::\text{int}, \lambda y. \text{twoCube2}(1, y)),$ 
(1,  $\lambda x. \text{twoCube2}(x, 0)), (-1, \lambda x. \text{twoCube2}(x, 1))\} = 4$ 
  using all-two-cubes-have-four-distict-edges that(2)
  by (auto simp add: boundary-def vertical-boundary-def horizontal-boundary-def)
  then show ?thesis
  by (auto simp: card-insert-if True split: if-split-asm)
  next
  case False
  then show ?thesis
  using no-shared-edges-have-similar-orientations
  by (simp add: that boundary-def vertical-boundary-def horizontal-boundary-def)
  qed
  then have  $\bigcup ((\lambda \text{twoCube}. \{(-1, \lambda y. \text{twoCube}(0, y)), (1, \lambda y. \text{twoCube}(1,$ 
y))\}) 'two-chain)
 $\cap \bigcup ((\lambda \text{twoCube}. \{(1::\text{int}, \lambda y. \text{twoCube}(y, 0)), (-1, \lambda y. \text{twoCube}(y,$ 
1))\}) 'two-chain)
= {}
  using Complete-Lattices.Union-disjoint by force
  then show ?thesis by force
  qed
qed
have one-chain-line-integral F {i} (two-chain-boundary two-chain)
= one-chain-line-integral F {i} (two-chain-vertical-boundary two-chain) +
one-chain-line-integral F {i} (two-chain-horizontal-boundary two-chain)
  using boundary-is-vert-hor horiz-verti-disjoint
  by (auto simp add: one-chain-line-integral-def hor-vert-finite sum.union-disjoint)

```

then have *x-axis-line-integral-is-only-horizontal*:
one-chain-line-integral $F \{i\}$ (*two-chain-boundary two-chain*)
 = *one-chain-line-integral* $F \{i\}$ (*two-chain-horizontal-boundary two-chain*)
using *vert-line-integral-zero* **by** *auto*

obtain $\mathcal{H} \mathcal{V}$ **where** *hv-props: finite* \mathcal{V}
 $(\forall (k, \gamma) \in \mathcal{V}. (\exists (k', \gamma') \in \text{two-chain-vertical-boundary two-chain}.$
 $(\exists a \in \{0 .. 1\}, \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma))) \ \gamma = \mathcal{H} \cup \mathcal{V}$
 (*common-sudiv-exists* (*two-chain-horizontal-boundary two-chain*) \mathcal{H})
 \vee *common-reparam-exists* \mathcal{H} (*two-chain-horizontal-boundary two-chain*)
finite \mathcal{H}
boundary-chain \mathcal{H}
 $\forall (k, \gamma) \in \mathcal{H}. \text{valid-path } \gamma$
using *only-vertical-division* **by** (*auto simp add: only-vertical-division-def*)
have *finite* $\{i\}$ **by** *auto*
then have *eq-integrals*: *one-chain-line-integral* $F \{i\}$ $\mathcal{H} =$ *one-chain-line-integral*
 $F \{i\}$ (*two-chain-horizontal-boundary two-chain*)
proof (*cases common-sudiv-exists* (*two-chain-horizontal-boundary two-chain*) \mathcal{H})
case *True*
then show *?thesis*
using *gen-common-subdivision-imp-eq-line-integral(1)* [*OF True two-chain-horizontal-boundary-is-boundary*
hv-props(6) - hor-vert-finite(2) hv-props(5)]
typeI-cube-line-integral-exists-boundary''
by *force*
next
case *False*
have *integ-exis-horiz*:
 $\forall (k, \gamma) \in \text{two-chain-horizontal-boundary two-chain}. \text{line-integral-exists } F \{i\}$
 γ
using *typeI-cube-line-integral-exists-boundary''* *assms*
by (*fastforce simp add: valid-two-chain-def*)
have *integ-exis*: $\forall (k, \gamma) \in \text{two-chain-boundary two-chain}. \text{line-integral-exists } F$
 $\{i\} \ \gamma$
using *typeI-cube-line-integral-exists-boundary* **by** *blast*
have *valid-paths*: $\forall (k, \gamma) \in \text{two-chain-vertical-boundary two-chain}. \text{valid-path } \gamma$
using *typeI-edges-are-valid-paths* *assms*
by (*fastforce simp add: two-chain-boundary-def two-chain-vertical-boundary-def*
boundary-def)
have *integ-exis-vert*:
 $(\bigwedge k \ \gamma. (\exists (k', \gamma') \in \text{two-chain-vertical-boundary two-chain}. \exists a \in \{0..1\}. \exists b \in \{0..1\}.$
 $a \leq b \wedge \text{subpath } a \ b \ \gamma' = \gamma) \implies$
line-integral-exists $F \{i\} \ \gamma)$
using *integ-exis valid-paths line-integral-exists-subpath* [*of* $F \{i\}$]
by (*fastforce simp add: two-chain-boundary-def two-chain-vertical-boundary-def*
two-chain-horizontal-boundary-def boundary-def)
have *finite-i*: *finite* $\{i\}$ **by** *auto*
have *i*: $\{i\} \subseteq \text{Basis}$ **using** *i-is-x-axis real-pair-basis* **by** *auto*
have *ii*: $\forall (k2, \gamma2) \in \text{two-chain-horizontal-boundary two-chain}. \forall b \in \{i\}. \text{contin-$
uous-on (*path-image* $\gamma2$) ($\lambda x. F \ x \cdot b$)

```

    using assms(1) field-cont-on-typeI-region-cont-on-edges assms(2)
  by (fastforce simp add: analytically-valid-def two-chain-horizontal-boundary-def
boundary-def path-image-def)
  have *: common-reparam-exists  $\mathcal{H}$  (two-chain-horizontal-boundary two-chain)
    using hv-props(4) False by auto
  show one-chain-line-integral  $F \{i\} \mathcal{H} = \text{one-chain-line-integral } F \{i\}$  (two-chain-horizontal-boundary
two-chain)
    using common-reparam-exists-imp-eq-line-integral(1)[OF finite-i hv-props(5)
hor-vert-finite(2) hv-props(6) two-chain-horizontal-boundary-is-boundary-chain ii
* hv-props(7) type-I-chain-horiz-bound-valid]
    by fastforce
qed

```

```

have line-integral-on-path:
  one-chain-line-integral  $F \{i\} \gamma =$ 
  one-chain-line-integral  $F \{i\}$  (two-chain-horizontal-boundary two-chain)
proof (auto simp add: one-chain-line-integral-def)
  have line-integral  $F \{i\}$  (snd oneCube) = 0 if oneCube: oneCube  $\in \mathcal{V}$  for
oneCube
  proof -
    obtain  $k \gamma$  where  $k\text{-gamma}: (k, \gamma) = \text{oneCube}$ 
      by (metis coeff-cube-to-path.cases)
    then obtain  $k' \gamma' a b$  where  $kp\text{-gammmap}$ :
      ( $k'::\text{int}, \gamma'$ )  $\in$  two-chain-vertical-boundary two-chain
       $a \in \{0 .. 1\}$ 
       $b \in \{0..1\}$ 
      subpath  $a b \gamma' = \gamma$ 
      using hv-props oneCube
      by (smt case-prodE split-conv)
    obtain  $x y1 y2$  where  $vert\text{-edge-def}: (k', \gamma') = (k', (\lambda t::\text{real}. (x::\text{real}, (1 - t)
* (y2) + t * y1)))$ 
      using valid-typeI-div  $kp\text{-gammmap}$ 
      by (auto simp add: typeI-twoCube-def two-chain-vertical-boundary-def verti-
cal-boundary-def)
    have  $vert\text{-edge-x-const}: \forall t. \gamma (t) \cdot i = x$ 
      by (metis (no-types, lifting) Pair-inject fstI i-is-x-axis inner-Pair-0(2)
 $kp\text{-gammmap}(4)$  real-inner-1-right subpath-def vert-edge-def)
    have  $\gamma = (\lambda t::\text{real}. (x::\text{real}, (1 - (b - a)*t - a) * (y2) + ((b-a)*t + a) *
y1))$ 
      using  $vert\text{-edge-def}$   $Product\text{-Type.snd-conv}$   $Product\text{-Type.fst-conv}$   $kp\text{-gammmap}(4)$ 
      by (simp add: subpath-def diff-diff-eq[symmetric])
    also have  $\dots = (\lambda t::\text{real}. (x::\text{real}, (1*y2 - a*y2) + a*y1 + ((b-a)*y1 - (b
- a)*y2)*t))$ 
      by (simp add: algebra-simps)
    finally have  $vert\text{-edge-is-straight-path}$ :
       $\gamma = (\lambda t::\text{real}. (x::\text{real}, (1*y2 - a*y2) + a*y1 + ((b-a)*y1 - (b -
a)*y2)*t))$  .
    show line-integral  $F \{i\}$  (snd oneCube) = 0
  proof -

```



```

have  $\forall x. \gamma$  differentiable at x
  by (simp add: straight-path-differentiable-x vert-edge-is-straight-path)
then have line-integral F {i}  $\gamma = 0$ 
  using line-integral-on-pair-straight-path(1) vert-edge-x-const by blast
then show ?thesis
  using Product-Type.snd-conv k-gamma by auto
qed
qed
then have  $\forall x \in \mathcal{V}. (\text{case } x \text{ of } (k, g) \Rightarrow (k::\text{int}) * \text{line-integral } F \{i\} g) = 0$ 
  by auto
then show  $(\sum x \in \gamma. \text{case } x \text{ of } (k, g) \Rightarrow \text{real-of-int } k * \text{line-integral } F \{i\} g) =$ 
 $(\sum x \in \text{two-chain-horizontal-boundary two-chain. case } x \text{ of } (k, g) \Rightarrow$ 
of-int k * line-integral F {i} g)
  using hv-props(1) hv-props(3) hv-props(5) sum-zero-set hor-vert-finite(2)
eq-integrals
  apply(auto simp add: one-chain-line-integral-def)
  by (smt Un-commute sum-zero-set)
qed
then have one-chain-line-integral F {i}  $\gamma =$ 
one-chain-line-integral F {i} (two-chain-boundary two-chain)
  using x-axis-line-integral-is-only-horizontal line-integral-on-path by auto
then show ?thesis
  using assms GreenThm-typeI-divisible by auto
qed
end

locale green-typeI-typeII-chain = R2: R2 i j + T1: green-typeI-chain i j F two-chain-typeI
+ T2: green-typeII-chain i j F two-chain-typeII for i j F two-chain-typeI two-chain-typeII
begin

lemma GreenThm-typeI-typeII-divisible-region-boundary:
assumes
  gen-divisions: gen-division s (cubeImage ' two-chain-typeI)
  gen-division s (cubeImage ' two-chain-typeII) and
  typeI-two-cubes-trace-horizontal-boundaries:
  two-chain-horizontal-boundary two-chain-typeI  $\subseteq \gamma$  and
  typeII-two-cubes-trace-vertical-boundaries:
  two-chain-vertical-boundary two-chain-typeII  $\subseteq \gamma$  and
  boundary-of-region-is-subset-of-partition-boundaries:
   $\gamma \subseteq \text{two-chain-boundary two-chain-typeI}$ 
   $\gamma \subseteq \text{two-chain-boundary two-chain-typeII}$ 
shows integral s ( $\lambda x. \text{partial-vector-derivative } (\lambda a. F a \cdot j) i x - \text{partial-vector-derivative}$ 
( $\lambda a. F a \cdot i) j x$ )
   $= \text{one-chain-line-integral } F \{i, j\} \gamma$ 
proof –
  let ?F-b' = partial-vector-derivative ( $\lambda a. F a \cdot i) j$ 
  let ?F-a' = partial-vector-derivative ( $\lambda a. F a \cdot j) i$ 
  have typeI-regions-integral: integral s ( $\lambda x. - \text{partial-vector-derivative } (\lambda a. F a \cdot$ 

```

i) $j x) = \text{one-chain-line-integral } F \{i\} \gamma$
using *T1.GreenThm-typeI-divisible-region-boundary*
gen-divisions(1) typeI-two-cubes-trace-horizontal-boundaries
boundary-of-region-is-subset-of-partition-boundaries(1)
by *blast*
have *typeII-regions-integral: integral s (partial-vector-derivative ($\lambda x. F x \cdot j$) i)*
= one-chain-line-integral F {j} γ
using *T2.GreenThm-typeII-divisible-region-boundary gen-divisions(2)*
typeII-two-cubes-trace-vertical-boundaries
boundary-of-region-is-subset-of-partition-boundaries(2)
by *auto*
have *integral-dis: integral s ($\lambda x. ?F-a' x - ?F-b' x$) = integral s ($\lambda x. ?F-a' x$) +*
integral s ($\lambda x. - ?F-b' x$)
proof *-*
have *$\forall \text{twoCube} \in \text{two-chain-typeII}. (?F-a' \text{ has-integral integral (cubeImage$*
twoCube) ?F-a') (cubeImage twoCube)
by *(simp add: analytically-valid-imp-part-deriv-integrable-on T2.F-anal-valid*
has-integral-iff)
then have *$\bigwedge u. u \in (\text{cubeImage ' two-chain-typeII}) \implies (?F-a' \text{ has-integral$*
integral u ?F-a') u
by *auto*
then have *(?F-a' has-integral ($\sum \text{img} \in \text{cubeImage ' two-chain-typeII}. \text{integral$*
img ?F-a')) s
using *gen-divisions(2) unfolding gen-division-def*
by *(metis has-integral-Union)*
then have *F-a'-integrable: (?F-a' integrable-on s) by auto*
have *$\forall \text{twoCube} \in \text{two-chain-typeI}. (?F-b' \text{ has-integral integral (cubeImage$*
twoCube) ?F-b') (cubeImage twoCube)
using *analytically-valid-imp-part-deriv-integrable-on T1.F-anal-valid by blast*
then have *$\bigwedge u. u \in (\text{cubeImage ' two-chain-typeI}) \implies (?F-b' \text{ has-integral$*
integral u ?F-b') u
by *auto*
then have *(?F-b' has-integral ($\sum \text{img} \in \text{cubeImage ' two-chain-typeI}. \text{integral$*
img ?F-b')) s
using *gen-divisions(1) unfolding gen-division-def*
by *(metis has-integral-Union)*
then show *?thesis*
by *(simp add: F-a'-integrable Henstock-Kurzweil-Integration.integral-diff has-integral-iff)*
qed
have *line-integral-dist: one-chain-line-integral F {i, j} γ = one-chain-line-integral*
F {i} γ + one-chain-line-integral F {j} γ
proof *(simp add: one-chain-line-integral-def)*
have *$k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral}$*
F {j} g
if *kg: (k,g) $\in \gamma$ for k g*
proof *-*
obtain *twoCube-typeI where twoCube-typeI-props:*
twoCube-typeI \in two-chain-typeI
(k, g) \in boundary twoCube-typeI

```

    typeI-twoCube twoCube-typeI
    continuous-on (cubeImage twoCube-typeI) (λx. F(x) · i)
using boundary-of-region-is-subset-of-partition-boundaries(1) two-chain-boundary-def
T1.valid-typeI-div
    T1.F-anal-valid kg
by (auto simp add: analytically-valid-def)
obtain twoCube-typeII where twoCube-typeII-props:
    twoCube-typeII ∈ two-chain-typeII
    (k, g) ∈ boundary twoCube-typeII
    typeII-twoCube twoCube-typeII
    continuous-on (cubeImage twoCube-typeII) (λx. F(x) · j)
using boundary-of-region-is-subset-of-partition-boundaries(2) two-chain-boundary-def
T2.valid-typeII-div
    kg T2.F-anal-valid
by (auto simp add: analytically-valid-def)
have line-integral F {i, j} g = line-integral F {i} g + line-integral F {j} g
proof –
    have int-exists-i: line-integral-exists F {i} g
        using T1.typeI-cube-line-integral-exists-boundary assms kg
        by (auto simp add: valid-two-chain-def)
    have int-exists-j: line-integral-exists F {j} g
        using T2.typeII-cube-line-integral-exists-boundary assms kg
        by (auto simp add: valid-two-chain-def)
    have finite: finite {i, j} by auto
    show ?thesis
        using line-integral-sum-gen[OF finite int-exists-i int-exists-j] R2.i-is-x-axis
R2.j-is-y-axis
    by auto
qed
    then show k * line-integral F {i, j} g = k * line-integral F {i} g + k *
line-integral F {j} g
    by (simp add: distrib-left)
qed
then have line-integral-distrib:
    (∑ (k,g)∈γ. k * line-integral F {i, j} g) =
    (∑ (k,g)∈γ. k * line-integral F {i} g + k * line-integral F {j} g)
    by (force intro: sum.cong split-cong)
    have (λx. (case x of (k, g) ⇒ (k::int) * line-integral F {i} g) + (case x of (k,
g) ⇒ (k::int) * line-integral F {j} g)) =
    (λx. (case x of (k, g) ⇒ (k * line-integral F {i} g) +
(k::int) * line-integral F {j} g))
    using comm-monoid-add-class.sum.distrib by auto
    then show (∑ (k, g)∈γ. k * line-integral F {i, j} g) =
    (∑ (k, g)∈γ. (k::int) * line-integral F {i} g) + (∑ (k, g)∈γ. (k::int) *
line-integral F {j} g)
    using comm-monoid-add-class.sum.distrib[of (λ(k, g). k * line-integral F {i}
g) (λ(k, g). k * line-integral F {j} g) γ]
    line-integral-distrib
    by presburger

```

qed
show *?thesis*
 using *integral-dis line-integral-dist typeI-regions-integral typeII-regions-integral*
 by *auto*
qed

lemma *GreenThm-typeI-typeII-divisible-region'*:

assumes
 only-vertical-division:
 only-vertical-division one-chain-typeI two-chain-typeI
 boundary-chain one-chain-typeI and
 only-horizontal-division:
 only-horizontal-division one-chain-typeII two-chain-typeII
 boundary-chain one-chain-typeII and
 typeI-and-typeII-one-chains-have-gen-common-subdiv:
 common-sudiv-exists one-chain-typeI one-chain-typeII
shows *integral s (λx. partial-vector-derivative (λx. (F x) · j) i x - partial-vector-derivative*
(λx. (F x) · i) j x) = one-chain-line-integral F {i, j} one-chain-typeI
 integral s (λx. partial-vector-derivative (λx. (F x) · j) i x - partial-vector-derivative
(λx. (F x) · i) j x) = one-chain-line-integral F {i, j} one-chain-typeII

proof –

let *?F-b' = partial-vector-derivative (λx. (F x) · i) j*
let *?F-a' = partial-vector-derivative (λx. (F x) · j) i*
have *one-chain-i-integrals:*
 one-chain-line-integral F {i} one-chain-typeI = one-chain-line-integral F {i}
one-chain-typeII ∧
 (∀ (k, γ) ∈ one-chain-typeI. line-integral-exists F {i} γ) ∧
 (∀ (k, γ) ∈ one-chain-typeII. line-integral-exists F {i} γ)

proof (*intro conjI*)

have *finite two-chain-typeI*
 using *T1.valid-typeI-div finite-image-iff*
 by (*auto simp add: gen-division-def valid-two-chain-def*)
then show *ii: ∀ (k, γ) ∈ one-chain-typeI. line-integral-exists F {i} γ*
 using *T1.members-of-only-vertical-div-line-integrable' assms*
 by *fastforce*
have *finite (two-chain-horizontal-boundary two-chain-typeI)*
by (*meson T1.valid-typeI-div finite-imageD finite-two-chain-horizontal-boundary*
gen-division-def valid-two-chain-def)
then have *finite one-chain-typeI*
 using *only-vertical-division(1) only-vertical-division-def by auto*
moreover have *finite one-chain-typeII*
 using *only-horizontal-division(1) only-horizontal-division-def by auto*
ultimately show *one-chain-line-integral F {i} one-chain-typeI = one-chain-line-integral*
F {i} one-chain-typeII
 and *∀ (k, γ) ∈ one-chain-typeII. line-integral-exists F {i} γ*
 using *gen-common-subdivision-imp-eq-line-integral[OF typeI-and-typeII-one-chains-have-gen-common-subdiv*
 only-vertical-division(2) only-horizontal-division(2)] ii
 by *auto*

qed

have *one-chain-j-integrals*:
 $one-chain-line-integral\ F\ \{j\}\ one-chain-typeII = one-chain-line-integral\ F\ \{j\}$
 $one-chain-typeI \wedge$
 $(\forall (k,\gamma) \in one-chain-typeII. line-integral-exists\ F\ \{j\}\ \gamma) \wedge$
 $(\forall (k,\gamma) \in one-chain-typeI. line-integral-exists\ F\ \{j\}\ \gamma)$

proof (*intro conjI*)
have *finite two-chain-typeII*
using *T2.valid-typeII-div finite-image-iff*
by (*auto simp add: gen-division-def valid-two-chain-def*)
then show *ii: $\forall (k,\gamma) \in one-chain-typeII. line-integral-exists\ F\ \{j\}\ \gamma$*
using *T2.members-of-only-horiz-div-line-integrable' assms T2.two-chain-valid-valid-cubes*
by *blast*

have *typeII-and-typI-one-chains-have-common-subdiv: common-sudiv-exists one-chain-typeII*
 $one-chain-typeI$
by (*simp add: common-sudiv-exists-comm typeI-and-typII-one-chains-have-gen-common-subdiv*)
have *iv: finite one-chain-typeI*
using *only-vertical-division(1) only-vertical-division-def by auto*
moreover have *iv': finite one-chain-typeII*
using *only-horizontal-division(1) only-horizontal-division-def by auto*
ultimately show $one-chain-line-integral\ F\ \{j\}\ one-chain-typeII =$
 $one-chain-line-integral\ F\ \{j\}\ one-chain-typeI$
 $\forall (k, \gamma) \in one-chain-typeI. line-integral-exists\ F\ \{j\}\ \gamma$
using *gen-common-subdivision-imp-eq-line-integral[OF typeII-and-typI-one-chains-have-common-subdiv*
 $only-horizontal-division(2)\ only-vertical-division(2)\ ii]\ ii$
by *auto*

qed

have *typeI-regions-integral*:
 $integral\ s\ (\lambda x. -\ ?F-b'\ x) = one-chain-line-integral\ F\ \{i\}\ one-chain-typeI$
using *T1.GreenThm-typeI-divisible-region-boundary-gen T1.valid-typeI-div*
 $T1.F-anal-valid\ only-vertical-division(1)$
by *auto*

have *typeII-regions-integral*:
 $integral\ s\ ?F-a' = one-chain-line-integral\ F\ \{j\}\ one-chain-typeII$
using *T2.GreenThm-typeII-divisible-region-boundary-gen T2.valid-typeII-div*
 $T2.F-anal-valid\ only-horizontal-division(1)$
by *auto*

have *line-integral-dist*:
 $one-chain-line-integral\ F\ \{i, j\}\ one-chain-typeI = one-chain-line-integral\ F\ \{i\}$
 $one-chain-typeI + one-chain-line-integral\ F\ \{j\}\ one-chain-typeI \wedge$
 $one-chain-line-integral\ F\ \{i, j\}\ one-chain-typeII = one-chain-line-integral\ F$
 $\{i\}\ one-chain-typeII + one-chain-line-integral\ F\ \{j\}\ one-chain-typeII$
proof (*simp add: one-chain-line-integral-def*)
have *line-integral-distrib*:
 $(\sum (k,g) \in one-chain-typeI. k * line-integral\ F\ \{i, j\}\ g) =$
 $(\sum (k,g) \in one-chain-typeI. k * line-integral\ F\ \{i\}\ g + k * line-integral\ F\ \{j\}$
 $g) \wedge$
 $(\sum (k,g) \in one-chain-typeII. k * line-integral\ F\ \{i, j\}\ g) =$
 $(\sum (k,g) \in one-chain-typeII. k * line-integral\ F\ \{i\}\ g + k * line-integral\ F$
 $\{j\}\ g)$

proof –
have $0: k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$
if $(k, g) \in \text{one-chain-typeII}$ **for** $k g$
proof –
have $\text{line-integral-exists } F \{i\} g \text{ line-integral-exists } F \{j\} g \text{ finite } \{i, j\}$
using $\text{one-chain-i-integrals one-chain-j-integrals that by fastforce+}$
moreover have $\{i\} \cap \{j\} = \{\}$
by $(\text{simp add: } R2.i\text{-is-x-axis } R2.j\text{-is-y-axis})$
ultimately have $\text{line-integral } F \{i, j\} g = \text{line-integral } F \{i\} g + \text{line-integral } F \{j\} g$
by $(\text{metis insert-is-Un line-integral-sum-gen}(1))$
then show $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$
by $(\text{simp add: distrib-left})$
qed
have $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$
if $(k, g) \in \text{one-chain-typeI}$ **for** $k g$
proof –
have $\text{line-integral } F \{i, j\} g = \text{line-integral } F \{i\} g + \text{line-integral } F \{j\} g$
by $(\text{smt that disjoint-insert}(2) \text{ finite.emptyI finite.insertI } R2.i\text{-is-x-axis inf-bot-right insert-absorb insert-commute insert-is-Un } R2.j\text{-is-y-axis line-integral-sum-gen}(1) \text{ one-chain-i-integrals one-chain-j-integrals prod.case-eq-if singleton-inject snd-conv})$
then show $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$
by $(\text{simp add: distrib-left})$
qed
then show $?thesis$
using 0 **by** $(\text{smt sum.cong split-cong})$
qed
show $(\sum_{(k::int, g) \in \text{one-chain-typeI}} k * \text{line-integral } F \{i, j\} g) =$
 $(\sum_{(k, g) \in \text{one-chain-typeI}} k * \text{line-integral } F \{i\} g) + (\sum_{(k::int, g) \in \text{one-chain-typeI}} k * \text{line-integral } F \{j\} g) \wedge$
 $(\sum_{(k::int, g) \in \text{one-chain-typeII}} k * \text{line-integral } F \{i, j\} g) =$
 $(\sum_{(k, g) \in \text{one-chain-typeII}} k * \text{line-integral } F \{i\} g) + (\sum_{(k::int, g) \in \text{one-chain-typeII}} k * \text{line-integral } F \{j\} g)$
proof –
have $0: (\lambda x. (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{i\} g) + (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{j\} g)) =$
 $(\lambda x. (\text{case } x \text{ of } (k::int, g) \Rightarrow (k * \text{line-integral } F \{i\} g) + k * \text{line-integral } F \{j\} g))$
using $\text{comm-monoid-add-class.sum.distrib}$ **by** auto
then have $1: (\sum_{x \in \text{one-chain-typeI}} (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{i\} g) + (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{j\} g)) =$
 $(\sum_{x \in \text{one-chain-typeI}} (\text{case } x \text{ of } (k::int, g) \Rightarrow (k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g)))$
by presburger
have $(\sum_{x \in \text{one-chain-typeII}} (\text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{i\} g))$

$+ (\text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{j\} g)) =$
 $(\sum x \in \text{one-chain-typeII}. (\text{case } x \text{ of } (k, g) \Rightarrow (k * \text{line-integral } F \{i\} g +$
 $k * \text{line-integral } F \{j\} g)))$
using 0 **by** *presburger*
then show *?thesis*
using *sum.distrib[of ($\lambda(k, g). k * \text{line-integral } F \{i\} g$) ($\lambda(k, g). k * \text{line-integral } F \{j\} g$) one-chain-typeI]*
*sum.distrib[of ($\lambda(k, g). k * \text{line-integral } F \{i\} g$) ($\lambda(k, g). k * \text{line-integral } F \{j\} g$) one-chain-typeII]*
line-integral-distrib 1
by *auto*
qed
qed
have *integral-dis: integral s ($\lambda x. ?F-a' x - ?F-b' x$) = integral s ($\lambda x. ?F-a' x$) + integral s ($\lambda x. - ?F-b' x$)*
proof –
have (*?F-a' has-integral integral (cubeImage twoCube) ?F-a'*) (*cubeImage twoCube*)
if *twoCube* \in *two-chain-typeII* **for** *twoCube*
by (*simp add: analytically-valid-imp-part-deriv-integrable-on T2.F-anal-valid has-integral-integrable-integral that*)
then have $\bigwedge u. u \in (\text{cubeImage } ' \text{two-chain-typeII}) \implies (?F-a' \text{ has-integral integral } u \text{ } ?F-a') u$
by *auto*
then have (*?F-a' has-integral ($\sum \text{img} \in \text{cubeImage } ' \text{two-chain-typeII}. \text{integral img } ?F-a')$) s*
using *T2.valid-typeII-div unfolding gen-division-def*
by (*metis has-integral-Union*)
then have *F-a'-integrable:*
(*?F-a' integrable-on s*) **by** *auto*
have $\forall \text{twoCube} \in \text{two-chain-typeI}. (?F-b' \text{ has-integral integral (cubeImage twoCube) } ?F-b') (\text{cubeImage twoCube})$
using *analytically-valid-imp-part-deriv-integrable-on T1.F-anal-valid* **by** *blast*
then have $\bigwedge u. u \in (\text{cubeImage } ' \text{two-chain-typeI}) \implies (?F-b' \text{ has-integral integral } u \text{ } ?F-b') u$
by *auto*
then have (*?F-b' has-integral ($\sum \text{img} \in \text{cubeImage } ' \text{two-chain-typeI}. \text{integral img } ?F-b')$) s*
using *T1.valid-typeI-div unfolding gen-division-def*
by (*metis has-integral-Union*)
then show *?thesis*
by (*simp add: F-a'-integrable Henstock-Kurzweil-Integration.integral-diff has-integral-iff*)
qed
show *integral s ($\lambda x. ?F-a' x - ?F-b' x$) = one-chain-line-integral F {i, j} one-chain-typeI*
using *one-chain-j-integrals integral-dis line-integral-dist typeI-regions-integral typeII-regions-integral*
by *auto*
show *integral s ($\lambda x. ?F-a' x - ?F-b' x$) = one-chain-line-integral F {i, j} one-chain-typeII*

using *one-chain-i-integrals integral-dis line-integral-dist typeI-regions-integral typeII-regions-integral*
by *auto*
qed

lemma *GreenThm-typeI-typeII-divisible-region:*

assumes *only-vertical-division:*
only-vertical-division one-chain-typeI two-chain-typeI
boundary-chain one-chain-typeI **and**
only-horizontal-division:
only-horizontal-division one-chain-typeII two-chain-typeII
boundary-chain one-chain-typeII **and**
typeI-and-typeII-one-chains-have-common-subdiv:
common-boundary-sudivision-exists one-chain-typeI one-chain-typeII
shows *integral s (λx. partial-vector-derivative (λx. (F x) · j) i x - partial-vector-derivative (λx. (F x) · i) j x) = one-chain-line-integral F {i, j} one-chain-typeI*
integral s (λx. partial-vector-derivative (λx. (F x) · j) i x - partial-vector-derivative (λx. (F x) · i) j x) = one-chain-line-integral F {i, j} one-chain-typeII
using *GreenThm-typeI-typeII-divisible-region' only-vertical-division only-horizontal-division common-subdiv-imp-gen-common-subdiv[OF typeI-and-typeII-one-chains-have-common-subdiv]*
by *auto*

lemma *GreenThm-typeI-typeII-divisible-region-finite-holes:*

assumes *valid-cube-boundary: ∀ (k,γ)∈boundary C. valid-path γ* **and**
only-vertical-division:
only-vertical-division (boundary C) two-chain-typeI **and**
only-horizontal-division:
only-horizontal-division (boundary C) two-chain-typeII **and**
s-is-oneCube: s = cubeImage C
shows *integral (cubeImage C) (λx. partial-vector-derivative (λx. F x · j) i x - partial-vector-derivative (λx. F x · i) j x) =*
one-chain-line-integral F {i, j} (boundary C)
using *GreenThm-typeI-typeII-divisible-region[OF only-vertical-division two-cube-boundary-is-boundary only-horizontal-division two-cube-boundary-is-boundary common-boundary-subdiv-exists-refl[OF assms(1)]] s-is-oneCube*
by *auto*

lemma *GreenThm-typeI-typeII-divisible-region-equivalent-boundary:*

assumes
gen-divisions: gen-division s (cubeImage ' two-chain-typeI)
gen-division s (cubeImage ' two-chain-typeII) **and**
typeI-two-cubes-trace-horizontal-boundaries:
two-chain-horizontal-boundary two-chain-typeI ⊆ one-chain-typeI **and**
typeII-two-cubes-trace-vertical-boundaries:
two-chain-vertical-boundary two-chain-typeII ⊆ one-chain-typeII **and**
boundary-of-region-is-subset-of-partition-boundaries:
one-chain-typeI ⊆ two-chain-boundary two-chain-typeI
one-chain-typeII ⊆ two-chain-boundary two-chain-typeII **and**
typeI-and-typeII-one-chains-have-common-subdiv:

common-boundary-sudivision-exists one-chain-typeI one-chain-typeII
shows *integral s* ($\lambda x. \text{partial-vector-derivative } (\lambda x. (F x) \cdot j) i x - \text{partial-vector-derivative } (\lambda x. (F x) \cdot i) j x$) = *one-chain-line-integral F {i, j} one-chain-typeI*
integral s ($\lambda x. \text{partial-vector-derivative } (\lambda x. (F x) \cdot j) i x - \text{partial-vector-derivative } (\lambda x. (F x) \cdot i) j x$) = *one-chain-line-integral F {i, j} one-chain-typeII*
proof –
let $?F\text{-}b' = \text{partial-vector-derivative } (\lambda x. (F x) \cdot i) j$
let $?F\text{-}a' = \text{partial-vector-derivative } (\lambda x. (F x) \cdot j) i$
have *one-chain-i-integrals*:
one-chain-line-integral F {i} one-chain-typeI = one-chain-line-integral F {i}
one-chain-typeII \wedge
 $(\forall (k, \gamma) \in \text{one-chain-typeI}. \text{line-integral-exists } F \{i\} \gamma) \wedge$
 $(\forall (k, \gamma) \in \text{one-chain-typeII}. \text{line-integral-exists } F \{i\} \gamma)$
proof (*intro conjI*)
have *i*: *boundary-chain one-chain-typeI*
using *two-chain-boundary-is-boundary-chain boundary-chain-def*
boundary-of-region-is-subset-of-partition-boundaries(1)
by *blast*
have *i'*: *boundary-chain one-chain-typeII*
using *two-chain-boundary-is-boundary-chain boundary-chain-def*
boundary-of-region-is-subset-of-partition-boundaries(2)
by *blast*
have $\bigwedge k \gamma. (k, \gamma) \in \text{one-chain-typeI} \implies \text{line-integral-exists } F \{i\} \gamma$
using *T1.typeI-cube-line-integral-exists-boundary assms*
by (*fastforce simp add: valid-two-chain-def*)
then show *ii*: $\forall (k, \gamma) \in \text{one-chain-typeI}. \text{line-integral-exists } F \{i\} \gamma$ **by** *auto*
have *finite* (*two-chain-boundary two-chain-typeI*)
unfolding *two-chain-boundary-def*
proof (*rule finite-UN-I*)
show *finite two-chain-typeI*
using *T1.valid-typeI-div finite-imageD gen-division-def valid-two-chain-def*
by *auto*
show $\bigwedge a. a \in \text{two-chain-typeI} \implies \text{finite } (\text{boundary } a)$
by (*simp add: boundary-def horizontal-boundary-def vertical-boundary-def*)
qed
then have *finite one-chain-typeI*
using *boundary-of-region-is-subset-of-partition-boundaries(1) finite-subset* **by**
fastforce
moreover have *finite* (*two-chain-boundary two-chain-typeII*)
unfolding *two-chain-boundary-def*
proof (*rule finite-UN-I*)
show *finite two-chain-typeII*
using *T2.valid-typeII-div finite-imageD gen-division-def valid-two-chain-def*
by *auto*
show $\bigwedge a. a \in \text{two-chain-typeII} \implies \text{finite } (\text{boundary } a)$
by (*simp add: boundary-def horizontal-boundary-def vertical-boundary-def*)
qed
then have *finite one-chain-typeII*
using *boundary-of-region-is-subset-of-partition-boundaries(2) finite-subset* **by**

fastforce
ultimately show $\text{one-chain-line-integral } F \{i\} \text{ one-chain-typeI} = \text{one-chain-line-integral}$
 $F \{i\} \text{ one-chain-typeII}$
 $\forall (k, \gamma) \in \text{one-chain-typeII}. \text{line-integral-exists } F \{i\} \gamma$
using *ii common-subdivision-imp-eq-line-integral*[*OF typeI-and-typII-one-chains-have-common-subdiv*
i i' ii]
by auto
qed
have *one-chain-j-integrals*:
 $\text{one-chain-line-integral } F \{j\} \text{ one-chain-typeI} = \text{one-chain-line-integral } F \{j\}$
 $\text{one-chain-typeII} \wedge$
 $(\forall (k, \gamma) \in \text{one-chain-typeI}. \text{line-integral-exists } F \{j\} \gamma) \wedge$
 $(\forall (k, \gamma) \in \text{one-chain-typeII}. \text{line-integral-exists } F \{j\} \gamma)$
proof (*intro conjI*)
have *i: boundary-chain one-chain-typeI and i': boundary-chain one-chain-typeII*
using *two-chain-boundary-is-boundary-chain boundary-of-region-is-subset-of-partition-boundaries*
unfolding *boundary-chain-def* **by** *blast+*
have *line-integral-exists F {j} γ if (k, γ) ∈ one-chain-typeII for k γ*
proof –
have *F-is-continuous: ∀ twoC ∈ two-chain-typeII. continuous-on (cubeImage*
twoC) (λa. F(a) · j)
using *T2.F-anal-valid* **by** (*simp add: analytically-valid-def*)
show *line-integral-exists F {j} γ*
using *that T2.valid-typeII-div*
boundary-of-region-is-subset-of-partition-boundaries(2)
using *green-typeII-cube.line-integral-exists-on-typeII-Cube-boundaries' assms*
valid-two-chain-def
apply (*simp add: two-chain-boundary-def*)
by (*metis T2.typeII-cube-line-integral-exists-boundary case-prodD subset-iff*
that two-chain-boundary-def)
qed
then show *ii: ∀ (k, γ) ∈ one-chain-typeII. line-integral-exists F {j} γ* **by auto**
have *finite (two-chain-boundary two-chain-typeI)*
unfolding *two-chain-boundary-def*
proof (*rule finite-UN-I*)
show *finite two-chain-typeI*
using *T1.valid-typeI-div finite-imageD gen-division-def valid-two-chain-def*
by auto
show $\bigwedge a. a \in \text{two-chain-typeI} \implies \text{finite (boundary } a)$
by (*simp add: boundary-def horizontal-boundary-def vertical-boundary-def*)
qed
then have *iv: finite one-chain-typeI*
using *boundary-of-region-is-subset-of-partition-boundaries(1) finite-subset*
by fastforce
have *finite (two-chain-boundary two-chain-typeII)*
unfolding *two-chain-boundary-def*
proof (*rule finite-UN-I*)
show *finite two-chain-typeII*
using *T2.valid-typeII-div finite-imageD gen-division-def valid-two-chain-def*

by *auto*
show $\bigwedge a. a \in \text{two-chain-typeII} \implies \text{finite (boundary } a)$
by (*simp add: boundary-def horizontal-boundary-def vertical-boundary-def*)
qed
then have *iv'*: *finite one-chain-typeII*
using *boundary-of-region-is-subset-of-partition-boundaries(2) finite-subset*
by *fastforce*
have *typeII-and-typI-one-chains-have-common-subdiv*:
common-boundary-sudivision-exists one-chain-typeII one-chain-typeI
using *typeI-and-typII-one-chains-have-common-subdiv*
common-boundary-sudivision-commutative
by *auto*
show *one-chain-line-integral* $F \{j\}$ *one-chain-typeI* = *one-chain-line-integral* F
 $\{j\}$ *one-chain-typeII*
 $\forall (k, \gamma) \in \text{one-chain-typeI}. \text{line-integral-exists } F \{j\} \gamma$
using *common-sudivision-imp-eq-line-integral[OF typeII-and-typI-one-chains-have-common-subdiv*
i' i ii iv' iw] ii
by *auto*
qed
have *typeI-regions-integral*:
integral $s (\lambda x. - ?F\text{-}b' x) = \text{one-chain-line-integral } F \{i\}$ *one-chain-typeI*
using *T1.GreenThm-typeI-divisible-region-boundary gen-divisions(1)*
typeI-two-cubes-trace-horizontal-boundaries
boundary-of-region-is-subset-of-partition-boundaries(1)
by *auto*
have *typeII-regions-integral*:
integral $s ?F\text{-}a' = \text{one-chain-line-integral } F \{j\}$ *one-chain-typeII*
using *T2.GreenThm-typeII-divisible-region-boundary gen-divisions(2)*
typeII-two-cubes-trace-vertical-boundaries
boundary-of-region-is-subset-of-partition-boundaries(2)
by *auto*
have *line-integral-dist*:
one-chain-line-integral $F \{i, j\}$ *one-chain-typeI* = *one-chain-line-integral* $F \{i\}$
one-chain-typeI + *one-chain-line-integral* $F \{j\}$ *one-chain-typeI* \wedge
one-chain-line-integral $F \{i, j\}$ *one-chain-typeII* = *one-chain-line-integral* F
 $\{i\}$ *one-chain-typeII* + *one-chain-line-integral* $F \{j\}$ *one-chain-typeII*
proof (*simp add: one-chain-line-integral-def*)
have *line-integral-distrib*:
 $(\sum (k,g) \in \text{one-chain-typeI}. k * \text{line-integral } F \{i, j\} g) =$
 $(\sum (k,g) \in \text{one-chain-typeI}. k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\}$
 $g) \wedge$
 $(\sum (k,g) \in \text{one-chain-typeII}. k * \text{line-integral } F \{i, j\} g) =$
 $(\sum (k,g) \in \text{one-chain-typeII}. k * \text{line-integral } F \{i\} g + k * \text{line-integral } F$
 $\{j\} g)$
proof –
have *0*: $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$
if $(k,g) \in \text{one-chain-typeII}$ **for** $k g$
proof –

```

have line-integral F {i, j} g = line-integral F {i} g + line-integral F {j} g
proof -
  have finite: finite {i, j} by auto
  have line-integral-all:  $\forall i \in \{i, j\}. \text{line-integral-exists } F \{i\} g$ 
    using one-chain-i-integrals one-chain-j-integrals that by auto
  show ?thesis
    using line-integral-sum-gen[OF finite] R2.i-is-x-axis R2.j-is-y-axis
line-integral-all by auto
qed
then show  $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$ 
  by (simp add: distrib-left)
qed
have  $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$ 
  if  $(k, g) \in \text{one-chain-typeI}$  for  $k g$ 
proof -
  have finite: finite {i, j} by auto
  have line-integral-all:  $\forall i \in \{i, j\}. \text{line-integral-exists } F \{i\} g$ 
    using one-chain-i-integrals one-chain-j-integrals that by auto
  have line-integral F {i, j} g = line-integral F {i} g + line-integral F {j} g
  using line-integral-sum-gen[OF finite] R2.i-is-x-axis R2.j-is-y-axis line-integral-all
by auto
then show  $k * \text{line-integral } F \{i, j\} g = k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g$ 
  by (simp add: distrib-left)
qed
then show ?thesis
  using 0 by (smt sum.cong split-cong)
qed
show  $(\sum_{(k::int, g) \in \text{one-chain-typeI}} k * \text{line-integral } F \{i, j\} g) =$ 
 $(\sum_{(k, g) \in \text{one-chain-typeI}} k * \text{line-integral } F \{i\} g) + (\sum_{(k::int, g) \in \text{one-chain-typeI}} k * \text{line-integral } F \{j\} g) \wedge$ 
 $(\sum_{(k::int, g) \in \text{one-chain-typeII}} k * \text{line-integral } F \{i, j\} g) =$ 
 $(\sum_{(k, g) \in \text{one-chain-typeII}} k * \text{line-integral } F \{i\} g) + (\sum_{(k::int, g) \in \text{one-chain-typeII}} k * \text{line-integral } F \{j\} g)$ 
proof -
  have 0:  $(\lambda x. (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{i\} g) + (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{j\} g)) =$ 
 $(\lambda x. (\text{case } x \text{ of } (k::int, g) \Rightarrow (k * \text{line-integral } F \{i\} g) + k * \text{line-integral } F \{j\} g))$ 
  using comm-monoid-add-class.sum.distrib by auto
  then have 1:  $(\sum_{x \in \text{one-chain-typeI}} (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{i\} g) + (\text{case } x \text{ of } (k::int, g) \Rightarrow k * \text{line-integral } F \{j\} g)) =$ 
 $(\sum_{x \in \text{one-chain-typeI}} (\text{case } x \text{ of } (k::int, g) \Rightarrow (k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g)))$ 
  by presburger
  have  $(\sum_{x \in \text{one-chain-typeII}} (\text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{i\} g) + (\text{case } x \text{ of } (k, g) \Rightarrow k * \text{line-integral } F \{j\} g)) =$ 

```

$(\sum x \in \text{one-chain-typeII}. (\text{case } x \text{ of } (k, g) \Rightarrow (k * \text{line-integral } F \{i\} g + k * \text{line-integral } F \{j\} g)))$
using 0 **by** presburger
then show ?thesis
using sum.distrib[of $(\lambda(k, g). k * \text{line-integral } F \{i\} g) (\lambda(k, g). k * \text{line-integral } F \{j\} g)$ one-chain-typeI]
sum.distrib[of $(\lambda(k, g). k * \text{line-integral } F \{i\} g) (\lambda(k, g). k * \text{line-integral } F \{j\} g)$ one-chain-typeII]
line-integral-distrib
1
by auto
qed
qed
have integral-dis: $\text{integral } s (\lambda x. ?F\text{-}a' x - ?F\text{-}b' x) = \text{integral } s (\lambda x. ?F\text{-}a' x) + \text{integral } s (\lambda x. - ?F\text{-}b' x)$
proof –
have (?F-a' has-integral $(\sum \text{img} \in \text{cubeImage 'two-chain-typeII}. \text{integral } \text{img } ?F\text{-}a')$) s
proof –
have (?F-a' has-integral $\text{integral } (\text{cubeImage } \text{twoCube}) ?F\text{-}a')$ (cubeImage twoCube)
if twoCube \in two-chain-typeII **for** twoCube
by (simp add: analytically-valid-imp-part-deriv-integrable-on T2.F-anal-valid has-integral-integrable-integral that)
then have $\bigwedge u. u \in (\text{cubeImage 'two-chain-typeII}) \implies (?F\text{-}a' \text{ has-integral } \text{integral } u ?F\text{-}a')$ u
by auto
then show ?thesis
using gen-divisions(2) **unfolding** gen-division-def
by (metis has-integral-Union)
qed
then have F-a'-integrable:
 $(?F\text{-}a' \text{ integrable-on } s)$ **by** auto
have (?F-b' has-integral $(\sum \text{img} \in \text{cubeImage 'two-chain-typeI}. \text{integral } \text{img } ?F\text{-}b')$) s
proof –
have $\forall \text{twoCube} \in \text{two-chain-typeI}. (?F\text{-}b' \text{ has-integral } \text{integral } (\text{cubeImage } \text{twoCube}) ?F\text{-}b')$ (cubeImage twoCube)
by (simp add: analytically-valid-imp-part-deriv-integrable-on T1.F-anal-valid has-integral-integrable-integral)
then have $\bigwedge u. u \in (\text{cubeImage 'two-chain-typeI}) \implies (?F\text{-}b' \text{ has-integral } \text{integral } u ?F\text{-}b')$ u
by auto
then show ?thesis
using gen-divisions(1) **unfolding** gen-division-def
by (metis has-integral-Union)
qed
then show ?thesis
using F-a'-integrable Henstock-Kurzweil-Integration.integral-diff **by** auto

```

qed
  show integral s ( $\lambda x. ?F\text{-}a' x - ?F\text{-}b' x$ ) = one-chain-line-integral F {i, j}
one-chain-typeI
  using one-chain-j-integrals integral-dis line-integral-dist typeI-regions-integral
typeII-regions-integral
  by auto
  show integral s ( $\lambda x. ?F\text{-}a' x - ?F\text{-}b' x$ ) = one-chain-line-integral F {i, j}
one-chain-typeII
  using one-chain-i-integrals integral-dis line-integral-dist typeI-regions-integral
typeII-regions-integral
  by auto
qed

end
end
theory SymmetricR2Shapes
  imports Green
begin

context R2
begin

lemma valid-path-valid-swap:
  assumes valid-path ( $\lambda x::real. ((f x)::real, (g x)::real)$ )
  shows valid-path (prod.swap o ( $\lambda x. (f x, g x)$ ))
  unfolding o-def valid-path-def piecewise-C1-differentiable-on-def swap-simp
proof (intro conjI)
  show continuous-on {0..1} ( $\lambda x. (g x, f x)$ )
    using assms
    using continuous-on-Pair continuous-on-componentwise[where  $f = (\lambda x. (f x,$ 
g x))]
    by (auto simp add: real-pair-basis valid-path-def piecewise-C1-differentiable-on-def)
  show  $\exists S. \text{finite } S \wedge (\lambda x. (g x, f x)) \text{ C1-differentiable-on } \{0..1\} - S$ 
  proof -
    obtain  $S$  where finite S and S: ( $\lambda x. (f x, g x)$ ) C1-differentiable-on {0..1} -
S
    using assms
    by (auto simp add: real-pair-basis valid-path-def piecewise-C1-differentiable-on-def)
  have  $0: f \text{ C1-differentiable-on } \{0..1\} - S$  using  $S$  assms
    using C1-diff-components-2[of (1,0) ( $\lambda x. (f x, g x)$ )]
    by (auto simp add: real-pair-basis algebra-simps)
  have  $1: g \text{ C1-differentiable-on } \{0..1\} - S$  using  $S$  assms
    using C1-diff-components-2 [of (0,1), OF -  $S$ ] real-pair-basis by fastforce
  have  $*$ : ( $\lambda x. (g x, f x)$ ) C1-differentiable-on {0..1} -  $S$ 
    using  $0\ 1$  C1-differentiable-on-components[where  $f = (\lambda x. (g x, f x))$ ]
    by (auto simp add: real-pair-basis valid-path-def piecewise-C1-differentiable-on-def)
  then show ?thesis using  $\langle \text{finite } S \rangle$  by auto
qed
qed

```

lemma *pair-fun-components*: $C = (\lambda x. (C\ x \cdot i, C\ x \cdot j))$
by (*simp add: i-is-x-axis inner-Pair-0 j-is-y-axis*)

lemma *swap-pair-fun*: $(\lambda y. \text{prod.swap } (C\ (y, 0))) = (\lambda x. (C\ (x, 0) \cdot j, C\ (x, 0) \cdot i))$
by (*simp add: prod.swap-def i-is-x-axis inner-Pair-0 j-is-y-axis*)

lemma *swap-pair-fun'*: $(\lambda y. \text{prod.swap } (C\ (y, 1))) = (\lambda x. (C\ (x, 1) \cdot j, C\ (x, 1) \cdot i))$
by (*simp add: prod.swap-def i-is-x-axis inner-Pair-0 j-is-y-axis*)

lemma *swap-pair-fun''*: $(\lambda y. \text{prod.swap } (C\ (0, y))) = (\lambda x. (C\ (0,x) \cdot j, C\ (0,x) \cdot i))$
by (*simp add: prod.swap-def i-is-x-axis inner-Pair-0 j-is-y-axis*)

lemma *swap-pair-fun'''*: $(\lambda y. \text{prod.swap } (C\ (1, y))) = (\lambda x. (C\ (1,x) \cdot j, C\ (1,x) \cdot i))$
by (*simp add: prod.swap-def i-is-x-axis inner-Pair-0 j-is-y-axis*)

lemma *swap-valid-boundaries*:
assumes $\forall (k,\gamma) \in \text{boundary } C. \text{ valid-path } \gamma$
assumes $(k,\gamma) \in \text{boundary } (\text{prod.swap } o\ C\ o\ \text{prod.swap})$
shows *valid-path* γ
using *assms*
valid-path-valid-swap[of $\lambda x. (\lambda x. C\ (x, 0))\ x \cdot i\ \lambda x. (\lambda x. C\ (x, 0))\ x \cdot j$]
pair-fun-components[of $(\lambda x. C\ (x, 0))$]
pair-fun-components[of $(\lambda y. C\ (y, 0))$]
valid-path-valid-swap[of $\lambda x. (\lambda y. C\ (y, 1))\ x \cdot i\ \lambda x. (\lambda y. C\ (y, 1))\ x \cdot j$]
pair-fun-components[of $(\lambda y. C\ (y, 1))$]
pair-fun-components[of $(\lambda x. C\ (x, 1))$]
valid-path-valid-swap[of $\lambda x. (\lambda y. C\ (1,y))\ x \cdot i\ \lambda x. (\lambda y. C\ (1,y))\ x \cdot j$]
pair-fun-components[of $(\lambda y. C\ (1,y))$]
pair-fun-components[of $(\lambda x. C\ (1,x))$]
valid-path-valid-swap[of $\lambda x. (\lambda y. C\ (0,y))\ x \cdot i\ \lambda x. (\lambda y. C\ (0,y))\ x \cdot j$]
pair-fun-components[of $(\lambda y. C\ (0,y))$]
pair-fun-components[of $(\lambda x. C\ (0,x))$]
by (*auto simp add: boundary-def horizontal-boundary-def vertical-boundary-def o-def real-pair-basis swap-pair-fun swap-pair-fun' swap-pair-fun'' swap-pair-fun'''*)

lemma *prod-comp-eq*:
assumes $f = \text{prod.swap } o\ g$
shows $\text{prod.swap } o\ f = g$
using *swap-comp-swap assms*
by *fastforce*

lemma *swap-typeI-is-typeII*:
assumes *typeI-twoCube* C
shows *typeII-twoCube* $(\text{prod.swap } o\ C\ o\ \text{prod.swap})$

proof (*simp add: typeI-twoCube-def typeII-twoCube-def*)
obtain $a\ b\ g1\ g2$ **where** $C: a < b$
 $(\forall x \in \{a..b\}. g2\ x \leq g1\ x)$
 $cubeImage\ C = \{(x, y). x \in \{a..b\} \wedge y \in \{g2\ x..g1\ x\}\}$
 $C = (\lambda(x, y). ((1 - x) * a + x * b, (1 - y) * g2\ ((1 - x) * a + x * b) + y$
 $*\ g1\ ((1 - x) * a + x * b)))$
 $g1$ *piecewise-C1-differentiable-on* $\{a..b\}$
 $g2$ *piecewise-C1-differentiable-on* $\{a..b\}$
using *typeI-cube-explicit-spec[OF assms]*
by *blast*
show $\exists a\ b. a < b \wedge$
 $(\exists g1\ g2. (\forall x \in \{a..b\}. g2\ x \leq g1\ x) \wedge$
 $prod.swap \circ C \circ prod.swap =$
 $(\lambda(y, x). ((1 - y) * g2\ ((1 - x) * a + x * b) + y * g1\ ((1 - x) *$
 $a + x * b), (1 - x) * a + x * b)) \wedge$
 $g1$ *piecewise-C1-differentiable-on* $\{a..b\} \wedge g2$ *piecewise-C1-differentiable-on*
 $\{a..b\})$
using C **by** (*fastforce simp add: prod.swap-def o-def*)
qed

lemma *valid-cube-valid-swap*:
assumes *valid-two-cube* C
shows *valid-two-cube* $(prod.swap \circ C \circ prod.swap)$
using *assms unfolding valid-two-cube-def boundary-def horizontal-boundary-def*
vertical-boundary-def
apply (*auto simp: card-insert-if split: if-split-asm*)
apply (*metis swap-swap*)
done

lemma *twoChainVertDiv-of-itself*:
assumes *finite* C
 $\forall (k, \gamma) \in (two-chain-boundary\ C). valid-path\ \gamma$
shows *only-vertical-division* $(two-chain-boundary\ C)\ C$
proof (*clarsimp simp add: only-vertical-division-def*)
show $\exists \mathcal{V}\ \mathcal{H}. finite\ \mathcal{H} \wedge finite\ \mathcal{V} \wedge$
 $(\forall x \in \mathcal{V}. case\ x\ of\ (k, \gamma) \Rightarrow \exists x \in two-chain-vertical-boundary\ C. case\ x$
of $(k', \gamma') \Rightarrow \exists a \in \{0..1\}. \exists b \in \{0..1\}. a \leq b \wedge subpath\ a\ b\ \gamma' = \gamma) \wedge$
 $(common-sudiv-exists\ (two-chain-horizontal-boundary\ C)\ \mathcal{H} \vee$
 $common-reparam-exists\ \mathcal{H}\ (two-chain-horizontal-boundary\ C)) \wedge$
 $boundary-chain\ \mathcal{H} \wedge two-chain-boundary\ C = \mathcal{V} \cup \mathcal{H} \wedge (\forall (k, \gamma) \in \mathcal{H}.$
 $valid-path\ \gamma)$
proof (*intro exI*)
let $?\mathcal{H} = two-chain-horizontal-boundary\ C$
have $0: \forall (k, \gamma) \in ?\mathcal{H}. valid-path\ \gamma$ **using** *assms(2)*
by (*auto simp add: two-chain-horizontal-boundary-def two-chain-boundary-def*
boundary-def)
have $\bigwedge a\ b. (a, b) \in two-chain-vertical-boundary\ C \implies$
 $\exists x \in two-chain-vertical-boundary\ C. case\ x\ of\ (k', \gamma') \Rightarrow \exists a \in \{0..1\}.$
 $\exists c \in \{0..1\}. a \leq c \wedge subpath\ a\ c\ \gamma' = b$

by (*metis (mono-tags, lifting) atLeastAtMost-iff case-prod-conv le-numeral-extra(1) order-refl subpath-trivial*)

moreover have *common-sudiv-exists ?H ?H*

using *gen-common-boundary-sudiv-exists-refl-twochain-boundary[OF 0 two-chain-horizontal-boundary-is-b] by auto*

moreover have *boundary-chain ?H*

using *two-chain-horizontal-boundary-is-boundary-chain by auto*

moreover have $\bigwedge a b. (a, b) \in \text{two-chain-boundary } C \implies (a, b) \notin ?\mathcal{H} \implies (a, b) \in \text{two-chain-vertical-boundary } C$

by (*auto simp add: two-chain-boundary-def two-chain-horizontal-boundary-def two-chain-vertical-boundary-def boundary-def*)

moreover have $\bigwedge a b. (a, b) \in \text{two-chain-vertical-boundary } C \implies (a, b) \in \text{two-chain-boundary } C$

$\bigwedge a b. (a, b) \in ?\mathcal{H} \implies (a, b) \in \text{two-chain-boundary } C$

by (*auto simp add: two-chain-boundary-def two-chain-horizontal-boundary-def two-chain-vertical-boundary-def boundary-def*)

moreover have $\bigwedge a b. (a, b) \in ?\mathcal{H} \implies \text{valid-path } b$

using *0 by blast*

ultimately show *finite ?H* \wedge

finite (two-chain-vertical-boundary C) \wedge

($\forall x \in \text{two-chain-vertical-boundary } C.$

case x of (k, γ) $\implies \exists x \in \text{two-chain-vertical-boundary } C. \text{ case x of (k', γ') $\implies \exists a \in \{0..1\}. \exists b \in \{0..1\}. a \leq b \wedge \text{subpath } a b \gamma' = \gamma) \wedge$$

(common-sudiv-exists ?H ?H \vee

common-reparam-exists ?H ?H) \wedge

boundary-chain ?H \wedge two-chain-boundary C = two-chain-vertical-boundary C \cup ?H \wedge ($\forall (k, \gamma) \in ?\mathcal{H}. \text{valid-path } \gamma$)

by (*auto simp add: finite-two-chain-horizontal-boundary[OF assms(1)] finite-two-chain-vertical-boundary[OF assms(1)]*)

qed

qed

end

definition *x-coord where* $x\text{-coord} \equiv (\lambda t::\text{real}. t - 1/2)$

lemma *x-coord-smooth: x-coord C1-differentiable-on {a..b}*

by (*simp add: x-coord-def*)

lemma *x-coord-bounds:*

assumes $(0::\text{real}) \leq x \leq 1$

shows $-1/2 \leq x\text{-coord } x \wedge x\text{-coord } x \leq 1/2$

using *assms by (auto simp add: x-coord-def)*

lemma *x-coord-img: x-coord \cdot $\{(0::\text{real})..1\} = \{-1/2 .. 1/2\}$*

by (*auto simp add: x-coord-def image-def algebra-simps*)

lemma *x-coord-back-img: finite ($\{0..1\} \cap x\text{-coord } - \cdot \{x::\text{real}\}$)*

by (simp add: finite-vimageI inj-on-def x-coord-def)

abbreviation $rot\text{-}x\ t1\ t2 \equiv (if\ (t1 - 1/2) \leq 0\ then\ (2 * t2 - 1) * t1 + 1/2 :: real\ else\ 2 * t2 - 2 * t1 * t2 + t1 - 1/2 :: real)$

lemma $rot\text{-}x\text{-}ivl$:

assumes $0 \leq x$

$x \leq 1$

$0 \leq y$

$y \leq 1$

shows $0 \leq rot\text{-}x\ x\ y \wedge rot\text{-}x\ x\ y \leq 1$

proof -

have $i: \bigwedge a::real. a \leq 0 \implies 0 \leq y \implies y \leq 1 \implies -1/2 < a \implies (a * (1 - 2*y) \leq 1/2)$

proof -

have $0: \bigwedge a::real. a \leq 0 \implies 0 \leq y \implies y \leq 1 \implies -1/2 < a \implies (-a \leq 1/2)$

by (sos ((($A < 0 * A < 1$) * $R < 1$) + ($R < 1 * (R < 1/4 * [2*a + 1]^2$))))

have $1: \bigwedge a. a \leq 0 \implies 0 \leq y \implies y \leq 1 \implies -1/2 < a \implies (a * (1 - 2*y) \leq -a)$

by (sos ((($A < 0 * A < 1$) * $R < 1$) + ((($A \leq 0 * (A < 1 * R < 1)$) * ($R < 2/3 * [1]^2$)) + ((($A \leq 0 * (A \leq 2 * R < 1)$) * ($R < 2/3 * [1]^2$)) + (($A \leq 0 * (A \leq 2 * (A < 0 * R < 1))$) * ($R < 2/3 * [1]^2$))))))

show $\bigwedge a::real. a \leq 0 \implies 0 \leq y \implies y \leq 1 \implies -1/2 < a \implies (a * (1 - 2*y) \leq 1/2)$ using 0 1 by force

qed

have $*$: $(x * 2 + y * 4 \leq 3 + x * (y * 4)) = ((x - 1) \leq 1/2 + (x - 1) * (y * 2))$

by (sos ((($A < 0 * R < 1$) + (($A \leq 0 * R < 1$) * ($R < 2 * [1]^2$)))) & ((($A < 0 * R < 1$) + (($A \leq 0 * R < 1$) * ($R < 1/2 * [1]^2$))))))

show ?thesis

using assms

apply (auto simp add: algebra-simps divide-simps linorder-class.not-le)

apply (sos ((($A < 0 * R < 1$) + ((($A \leq 2 * (A \leq 3 * R < 1)$) * ($R < 1 * [1]^2$)) + ((($A \leq 1 * R < 1$) * ($R < 1 * [1]^2$)) + (($A \leq 0 * (A \leq 1 * R < 1)$) * ($R < 2 * [1]^2$))))))

apply (sos ((($A < 0 * R < 1$) + ((($A \leq 2 * R < 1$) * ($R < 1 * [1]^2$)) + ((($A \leq 1 * (A \leq 3 * R < 1)$) * ($R < 1 * [1]^2$)) + (($A \leq 0 * (A \leq 2 * R < 1)$) * ($R < 2 * [1]^2$))))))

using i[of (x::real) - 1] affine-ineq

apply (fastforce simp: algebra-simps *)

done

qed

end

2 The Circle Example

theory *CircExample*

```

imports Green SymmetricR2Shapes

begin

locale circle = R2 +
  fixes d::real
  assumes d-gt-0: 0 < d
begin

definition circle-y where
  circle-y t = sqrt (1/4 - t * t)

definition circle-cube where
  circle-cube = (λ(x,y). ((x - 1/2) * d, (2 * y - 1) * d * sqrt (1/4 - (x - 1/2)*(x
-1/2))))

lemma circle-cube-nice:
  shows circle-cube = (λ(x,y). (d * x-coord x, (2 * y - 1) * d * circle-y (x-coord
x)))
  by (auto simp add: circle-cube-def circle-y-def x-coord-def)

definition rot-circle-cube where
  rot-circle-cube = prod.swap ∘ (circle-cube) ∘ prod.swap

abbreviation rot-y t1 t2 ≡ ((t1 - 1/2)/(2 * circle-y (x-coord (rot-x t1 t2))) + 1/2)::real

definition x-coord-inv (x::real) = (1/2) + x

lemma x-coord-inv-1: x-coord-inv (x-coord (x::real)) = x
  by (auto simp add: x-coord-inv-def x-coord-def)

lemma x-coord-inv-2: x-coord (x-coord-inv (x::real)) = x
  by (auto simp add: x-coord-inv-def x-coord-def)

definition circle-y-inv = circle-y

abbreviation rot-x'' (x::real) (y::real) ≡ (x-coord-inv ((2 * y - 1) * circle-y
(x-coord x)))

lemma circle-y-bounds:
  assumes -1/2 ≤ (x::real) ∧ x ≤ 1/2
  shows 0 ≤ circle-y x circle-y x ≤ 1/2
  unfolding circle-y-def real-sqrt-ge-0-iff
proof -
  show 0 ≤ 1/4 - x * x
  using assms
  by (sos (((A < 0 * R < 1) + ((A <= 0 * (A <= 1 * R < 1)) * (R < 1 * [1]^2))))
  show sqrt (1/4 - x * x) ≤ 1/2
  apply (rule real-le-lsqrt)

```

using *assms* by(*auto simp add: divide-simps algebra-simps*)
qed

lemma *circle-y-x-coord-bounds*:

assumes $0 \leq (x::real) \ x \leq 1$

shows $0 \leq \text{circle-}y \ (x\text{-coord } x) \wedge \text{circle-}y \ (x\text{-coord } x) \leq 1/2$

using *circle-y-bounds*[*OF x-coord-bounds*][*OF assms*] by *auto*

lemma *rot-x-ivl*:

assumes $(0::real) \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1$

shows $0 \leq \text{rot-}x'' \ x \ y \wedge \text{rot-}x'' \ x \ y \leq 1$

proof

have $\bigwedge a::real. 0 \leq a \wedge a \leq 1/2 \implies 0 \leq 1/2 + (2 * y - 1) * a$ using *assms*

by (*sos* ((($A < 0 * R < 1$) + ((($A \leq 4 * R < 1$) * ($R < 1/2 * [1]^2$))) + ((($A \leq 1 * (A \leq 5 * R < 1)$) * ($R < 1 * [1]^2$))) + (($A \leq 0 * (A \leq 4 * R < 1)$) * ($R < 1 * [1]^2$))))))

then show $0 \leq \text{rot-}x'' \ x \ y$

using *assms circle-y-x-coord-bounds* by(*auto simp add: x-coord-inv-def*)

have $\bigwedge a::real. 0 \leq a \wedge a \leq 1/2 \implies 1/2 + (2 * y - 1) * a \leq 1$ using *assms*

by (*sos* ((($A < 0 * R < 1$) + ((($A \leq 5 * R < 1$) * ($R < 1/2 * [1]^2$))) + ((($A \leq 1 * (A \leq 4 * R < 1)$) * ($R < 1 * [1]^2$))) + (($A \leq 0 * (A \leq 5 * R < 1)$) * ($R < 1 * [1]^2$))))))

then show $\text{rot-}x'' \ x \ y \leq 1$

using *assms circle-y-x-coord-bounds* by (*auto simp add: x-coord-inv-def*)

qed

abbreviation *rot-y''* ($x::real$) ($y::real$) $\equiv (x\text{-coord } x)/(2 * (\text{circle-}y \ (x\text{-coord } (\text{rot-}x'' \ x \ y)))) + 1/2$

lemma *rot-y-ivl*:

assumes $(0::real) \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1$

shows $0 \leq \text{rot-}y'' \ x \ y \wedge \text{rot-}y'' \ x \ y \leq 1$

proof

show $0 \leq \text{rot-}y'' \ x \ y$

proof(*cases* ($x\text{-coord } x$) < 0)

case *True*

have $i: \bigwedge a \ b::real. a < 0 \implies 0 \leq a + b \implies (0 \leq a/(2*(b)) + 1/2)$

by(*auto simp add: algebra-simps divide-simps*)

have $*$: $(1/2 - x) \leq \text{sqrt} \ (x * x + (1/4 + (x * (y * 4) + x * (x * (y * (y * 4)))))) - (x + (x * (x * (y * 4) + x * (y * (y * 4))))))$

apply (*rule real-le-rsqrt*)

using *assms apply* (*simp add: algebra-simps power2-eq-square mult-left-le-one-le*)

by (*sos* ((($A < 0 * R < 1$) + (($A \leq 0 * (A \leq 1 * (A \leq 2 * (A \leq 3 * R < 1)))$)) * ($R < 4 * [1]^2$))))

have *rw*: $|x - x * x| = x - x * x$ using *assms*

by (*sos* (($\&$ ((($A < 0 * A < 1$) * $R < 1$) + (($A \leq 0 * (A \leq 1 * (A < 1 * R < 1))) * ($R < 1 * [1]^2$)))) & ((($A < 0 * A < 1$) * $R < 1$) + (($A \leq 0 * (A \leq 1 * (A < 1 * R < 1))) * ($R < 1 * [1]^2$))))))$$

have $0 \leq x\text{-coord } x + (\text{circle-}y \ (x\text{-coord } (\text{rot-}x'' \ x \ y)))$

```

    using * apply (auto simp add: x-coord-inv-2)
    by (auto simp add: circle-y-def algebra-simps rw x-coord-def)
  then show ?thesis
    using True i by blast
next
case False
have i:  $\bigwedge a b::real. 0 \leq a \implies 0 \leq b \implies (0 \leq a/(2*b) + 1/2)$ 
  by (auto simp add: algebra-simps divide-simps)
have  $0 \leq \text{circle-y } (x\text{-coord } (x\text{-coord-inv } ((2 * y - 1) * \text{circle-y } (x\text{-coord } x))))$ 
proof -
  have rw:  $|x - x * x| = x - x * x$  using assms
  by (sos ((() & (((((A<0 * A<1) * R<1) + ((A<=0 * (A<=1 * (A<1 * R<1))) * (R<1 * [1]^2)))) & (((((A<0 * A<1) * R<1) + ((A<=0 * (A<=1 * (A<1 * R<1))) * (R<1 * [1]^2))))))))))
  have  $\bigwedge x. 0 \leq x \implies x \leq 1/2 \implies -1/2 \leq (2 * y - 1) * x$  using assms
  by (sos (((A<0 * R<1) + (((A<=4 * R<1) * (R<1/2 * [1]^2)) + (((A<=1 * (A<=5 * R<1)) * (R<1 * [1]^2)) + ((A<=0 * (A<=4 * R<1)) * (R<1 * [1]^2))))))))
  then have  $-1/2 \leq (2 * y - 1) * \text{circle-y } (x\text{-coord } x)$ 
    using circle-y-x-coord-bounds assms(1-2) by auto
  moreover
  have  $\bigwedge x. 0 \leq x \implies x \leq 1/2 \implies (2 * y - 1) * x \leq 1/2$  using assms
  by (sos (((A<0 * R<1) + (((A<=5 * R<1) * (R<1/2 * [1]^2)) + (((A<=1 * (A<=4 * R<1)) * (R<1 * [1]^2)) + ((A<=0 * (A<=5 * R<1)) * (R<1 * [1]^2))))))))
  then have  $(2 * y - 1) * \text{circle-y } (x\text{-coord } x) \leq 1/2$ 
    using circle-y-x-coord-bounds assms(1-2) by auto
  ultimately show  $0 \leq \text{circle-y } (x\text{-coord } (x\text{-coord-inv } ((2 * y - 1) * \text{circle-y } (x\text{-coord } x))))$ 
    by (simp add: circle-y-bounds(1) x-coord-inv-2)
qed
then show ?thesis
  using False by auto
qed
have i:  $\bigwedge a b::real. a < 0 \implies 0 \leq b \implies (a/(2*(b)) + 1/2) \leq 1$ 
  by (auto simp add: algebra-simps divide-simps)
show rot-y'' x y  $\leq 1$ 
proof (cases (x-coord x) < 0)
case True
  have i:  $\bigwedge a b::real. a < 0 \implies 0 \leq b \implies (a/(2*(b)) + 1/2) \leq 1$ 
    by (auto simp add: algebra-simps divide-simps)
  have  $\bigwedge x. 0 \leq x \implies x \leq 1/2 \implies -1/2 \leq (2 * y - 1) * x$  using assms
  by (sos (((A<0 * R<1) + (((A<=4 * R<1) * (R<1/2 * [1]^2)) + (((A<=1 * (A<=5 * R<1)) * (R<1 * [1]^2)) + ((A<=0 * (A<=4 * R<1)) * (R<1 * [1]^2))))))))
  then have  $-1/2 \leq (2 * y - 1) * \text{circle-y } (x\text{-coord } x)$ 
    using circle-y-x-coord-bounds assms(1-2) by auto
  moreover have  $\bigwedge x. 0 \leq x \implies x \leq 1/2 \implies (2 * y - 1) * x \leq 1/2$  using
  assms

```

```

    by (sos (((A<0 * R<1) + (((A<=5 * R<1) * (R<1/2 * [1]^2)) + (((A<=1
* (A<=4 * R<1)) * (R<1 * [1]^2)) + ((A<=0 * (A<=5 * R<1)) * (R<1 *
[1]^2)))))))
    then have (2 * y - 1) * circle-y (x-coord x) ≤ 1/2
    using circle-y-x-coord-bounds assms(1-2) by auto
    ultimately have 0 ≤ circle-y (x-coord (x-coord-inv ((2 * y - 1) * circle-y
(x-coord x)))
    by (simp add: circle-y-bounds(1) x-coord-inv-2)
    then show ?thesis
    by (simp add: True i)
next
case False
have i:  $\bigwedge a b :: \text{real}. 0 \leq a \implies a \leq b \implies (a/(2*b) + 1/2) \leq 1$ 
by (auto simp add: algebra-simps divide-simps)
have (x - 1/2) * (x - 1/2) ≤ (x * x + (1/4 + (x * (y * 4) + x * (x * (y *
(y * 4)))))) - (x + (x * (x * (y * 4)) + x * (y * (y * 4))))
using assms False
apply (auto simp add: x-coord-def)
by (sos (((A<0 * R<1) + (((A<=0 * (A<=1 * (A<=2 * R<1))) * (R<2 *
[1]^2)) + ((A<=0 * (A<=1 * (A<=2 * (A<=3 * R<1)))) * (R<2 * [1]^2))))))
    then have sqrt ((x - 1/2) * (x - 1/2)) ≤ sqrt (x * x + (1/4 + (x * (y *
4) + x * (x * (y * (y * 4)))))) - (x + (x * (x * (y * 4)) + x * (y * (y * 4))))
    using real-sqrt-le-mono by blast
    then have *: (x - 1/2) ≤ sqrt (x * x + (1/4 + (x * (y * 4) + x * (x * (y *
(y * 4)))))) - (x + (x * (x * (y * 4)) + x * (y * (y * 4))))
    using assms False by (auto simp add: x-coord-def)
    have rw:  $|x - x * x| = x - x * x$  using assms
    by (sos (( & (((((A<0 * A<1) * R<1) + ((A<=0 * (A<=1 * (A<1 *
R<1))) * (R<1 * [1]^2)))) & (((A<0 * A<1) * R<1) + ((A<=0 * (A<=1 *
(A<1 * R<1))) * (R<1 * [1]^2)))))))
    have x-coord x ≤ circle-y (x-coord (x-coord-inv ((2 * y - 1) * circle-y (x-coord
x))))
    using * unfolding x-coord-inv-2
    by (auto simp add: circle-y-def algebra-simps rw x-coord-def)
    then show ?thesis
    using False i by auto
qed
qed

lemma circle-eq-rot-circle:
  assumes  $0 \leq x \leq 1$   $0 \leq y \leq 1$ 
  shows (circle-cube (x, y)) = (rot-circle-cube (rot-y'' x y, rot-x'' x y))
proof
  have rw:  $|1/4 - x\text{-coord } x * x\text{-coord } x| = 1/4 - x\text{-coord } x * x\text{-coord } x$ 
  apply (rule abs-of-nonneg)
  using assms mult-left-le by (auto simp add: x-coord-def divide-simps alge-
bra-simps)
  show fst (circle-cube (x, y)) = fst (rot-circle-cube (rot-y'' x y, rot-x'' x y))
  using assms d-gt-0

```

apply(*simp add: circle-cube-nice rot-circle-cube-def x-coord-inv-2 circle-y-def algebra-simps rw*)
apply(*auto simp add: x-coord-def algebra-simps*)
by (sos (((((A<0 * A<1) * ((A<0 * A<1) * R<1)) + (([~ 4*d^2] * A=0) + (((A<=1 * (A<=2 * (A<=3 * R<1))) * (R<8 * [d]^2)) + ((A<=1 * (A<=2 * (A<=3 * (A<1 * R<1)))) * (R<8 * [d]^2)))))) & (((A<0 * A<1) * ((A<0 * A<1) * R<1)) + (([~ 4*d^2] * A=0) + (((A<=0 * (A<=2 * (A<=3 * R<1))) * (R<8 * [d]^2)) + ((A<=0 * (A<=2 * (A<=3 * (A<1 * R<1)))) * (R<8 * [d]^2))))))))))
show *snd (circle-cube (x, y)) = snd (rot-circle-cube (rot-y'' x y, rot-x'' x y))*
using *assms*
by(*auto simp add: circle-cube-def rot-circle-cube-def x-coord-inv-def circle-y-def x-coord-def*)
qed

lemma *rot-circle-eq-circle:*

assumes $0 \leq x \leq 1 \ 0 \leq y \leq 1$
shows $(\text{rot-circle-cube } (x, y)) = (\text{circle-cube } (\text{rot-x}'' y x, \text{rot-y}'' y x))$
proof
show *fst (rot-circle-cube (x, y)) = fst (circle-cube (rot-x'' y x, rot-y'' y x))*
using *assms*
by(*auto simp add: circle-cube-def rot-circle-cube-def x-coord-inv-def circle-y-def x-coord-def*)
have *rw: |1/4 - x-coord y * x-coord y| = 1/4 - x-coord y * x-coord y*
apply(*rule abs-of-nonneg*)
using *assms mult-left-le* **by** (*auto simp add: x-coord-def divide-simps algebra-simps*)
show *snd (rot-circle-cube (x, y)) = snd (circle-cube (rot-x'' y x, rot-y'' y x))*
using *assms d-gt-0*
apply(*simp add: circle-cube-nice rot-circle-cube-def x-coord-inv-2 circle-y-def algebra-simps rw*)
apply(*auto simp add: x-coord-def algebra-simps*)
by (sos ((((((A<0 * A<1) * ((A<0 * A<1) * R<1)) + (([~ 4*d^2] * A=0) + (((A<=0 * (A<=1 * (A<=3 * R<1))) * (R<8 * [d]^2)) + ((A<=0 * (A<=1 * (A<=3 * (A<1 * R<1)))) * (R<8 * [d]^2)))))) & (((A<0 * A<1) * ((A<0 * A<1) * R<1)) + (([~ 4*d^2] * A=0) + (((A<=0 * (A<=1 * (A<=2 * R<1))) * (R<8 * [d]^2)) + ((A<=0 * (A<=1 * (A<=2 * (A<1 * R<1)))) * (R<8 * [d]^2)))))))))))))
qed

lemma *rot-img-eq:*

assumes $0 < d$
shows $(\text{cubeImage } (\text{circle-cube })) = (\text{cubeImage } (\text{rot-circle-cube}))$
apply(*auto simp add: cubeImage-def image-def cbox-def real-pair-basis*)
by (*meson rot-y-ivl rot-x-ivl assms circle-eq-rot-circle rot-circle-eq-circle*)+

lemma *rot-circle-div-circle:*

assumes $0 < (d::\text{real})$
shows *gen-division (cubeImage circle-cube) (cubeImage ‘ {rot-circle-cube})*

```

using rot-img-eq[OF assms] by(auto simp add: gen-division-def)

lemma circle-cube-boundary-valid:
  assumes  $(k,\gamma)\in\text{boundary circle-cube}$ 
  shows valid-path  $\gamma$ 
proof -
  have  $f01: \text{finite}\{0,1\}$ 
  by simp
  show ?thesis
  using assms
  unfolding boundary-def horizontal-boundary-def vertical-boundary-def circle-cube-def
valid-path-def piecewise-C1-differentiable-on-def
  by safe (rule derivative-intros continuous-intros f01 exI ballI conjI refl | force
simp add: field-simps)
qed

lemma rot-circle-cube-boundary-valid:
  assumes  $(k,\gamma)\in\text{boundary rot-circle-cube}$ 
  shows valid-path  $\gamma$ 
  using assms swap-valid-boundaries circle-cube-boundary-valid
  by (fastforce simp add: rot-circle-cube-def)

lemma diff-divide-cancel:
  fixes  $z::\text{real}$  shows  $z \neq 0 \implies (a * z - a * (b * z)) / z = (a - a * b)$ 
  by (auto simp: field-simps)

lemma circle-cube-is-type-I:
  assumes  $0 < d$ 
  shows typeI-twoCube circle-cube
  unfolding typeI-twoCube-def
proof (intro exI conjI ballI)
  have  $f01: \text{finite}\{-d/2,d/2\}$ 
  by simp
  show  $-d/2 < d/2$ 
  using assms by simp
  show  $(\lambda x. d * \text{sqrt}(1/4 - (x/d) * (x/d))) \text{ piecewise-C1-differentiable-on } \{-d/2..d/2\}$ 
  using assms unfolding piecewise-C1-differentiable-on-def
  apply (intro exI conjI)
  apply (rule ballI refl f01 derivative-intros continuous-intros | simp)
  apply (auto simp: field-simps)
  by sos
  show  $(\lambda x. -d * \text{sqrt}(1/4 - (x/d) * (x/d))) \text{ piecewise-C1-differentiable-on } \{-d/2..d/2\}$ 
  using assms unfolding piecewise-C1-differentiable-on-def
  apply (intro exI conjI)
  apply (rule ballI refl f01 derivative-intros continuous-intros | simp)
  apply (auto simp: field-simps)
  by sos

```



```

show  $- d * \text{sqrt} (1/4 - x / d * (x / d)) \leq d * \text{sqrt} (1/4 - x / d * (x / d))$ 
if  $x \in \{- d/2..d/2\}$  for  $x$ 
proof -
  have  $*$ :  $x^2 \leq (d/2)^2$ 
    using real-sqrt-le-iff that by fastforce
  show ?thesis
  apply (rule mult-right-mono)
  using assms * apply (simp-all add: divide-simps power2-eq-square)
  done
qed
qed (auto simp add: circle-cube-def divide-simps algebra-simps diff-divide-cancel)

```

```

lemma rot-circle-cube-is-type-II:
  shows typeII-twoCube rot-circle-cube
  using d-gt-0 swap-typeI-is-typeII circle-cube-is-type-I
  by (auto simp add: rot-circle-cube-def)

```

```

definition circle-bot-edge where
  circle-bot-edge = (1::int,  $\lambda t. (x\text{-coord } t * d, - d * \text{circle-y } (x\text{-coord } t))$ )

```

```

definition circle-top-edge where
  circle-top-edge = (- 1::int,  $\lambda t. (x\text{-coord } t * d, d * \text{circle-y } (x\text{-coord } t))$ )

```

```

definition circle-right-edge where
  circle-right-edge = (1::int,  $\lambda y. (d/2, 0)$ )

```

```

definition circle-left-edge where
  circle-left-edge = (- 1::int,  $\lambda y. (- (d/2), 0)$ )

```

```

lemma circle-cube-boundary-explicit:
  boundary circle-cube = {circle-left-edge, circle-right-edge, circle-bot-edge, circle-top-edge}
  by (auto simp add: valid-two-cube-def boundary-def horizontal-boundary-def vertical-boundary-def circle-cube-def
    circle-top-edge-def circle-bot-edge-def circle-cube-nice x-coord-def circle-y-def
    circle-left-edge-def circle-right-edge-def)

```

```

definition rot-circle-right-edge where
  rot-circle-right-edge = (1::int,  $\lambda t. (d * \text{circle-y } (x\text{-coord } t), x\text{-coord } t * d)$ )

```

```

definition rot-circle-left-edge where
  rot-circle-left-edge = (- 1::int,  $\lambda t. (- d * \text{circle-y } (x\text{-coord } t), x\text{-coord } t * d)$ )

```

```

definition rot-circle-top-edge where
  rot-circle-top-edge = (- 1::int,  $\lambda y. (0, d/2)$ )

```

```

definition rot-circle-bot-edge where
  rot-circle-bot-edge = (1::int,  $\lambda y. (0, - (d/2))$ )

```

```

lemma rot-circle-cube-boundary-explicit:

```

$boundary\ (rot\ circle\ cube) =$
 $\{rot\ circle\ top\ edge, rot\ circle\ bot\ edge, rot\ circle\ right\ edge, rot\ circle\ left\ edge\}$
by (*auto simp add: rot-circle-cube-def valid-two-cube-def boundary-def horizontal-boundary-def vertical-boundary-def circle-cube-def*
rot-circle-right-edge-def rot-circle-left-edge-def x-coord-def circle-y-def rot-circle-top-edge-def
rot-circle-bot-edge-def)

lemma *rot-circle-cube-vertical-boundary-explicit*:
 $vertical\ boundary\ rot\ circle\ cube = \{rot\ circle\ right\ edge, rot\ circle\ left\ edge\}$
by (*auto simp add: rot-circle-cube-def valid-two-cube-def vertical-boundary-def*
circle-cube-def
rot-circle-right-edge-def rot-circle-left-edge-def x-coord-def circle-y-def)

lemma *circ-left-edge-neq-top*:
 $(-1::int, \lambda y::real. -(d/2), 0) \neq (-1, \lambda x. ((x - 1/2) * d, d * sqrt(1/4 - (x - 1/2) * (x - 1/2))))$
by (*metis (no-types, lifting) add-diff-cancel-right' d-gt-0 mult.commute mult-cancel-left*
order-less-irrefl prod.inject)

lemma *circle-cube-valid-two-cube: valid-two-cube (circle-cube)*

proof (*auto simp add: valid-two-cube-def boundary-def horizontal-boundary-def vertical-boundary-def circle-cube-def*)

have *iv*: $(-1::int, \lambda y::real. -(d/2), 0) \neq (-1, \lambda x. ((x - 1/2) * d, d * sqrt(1/4 - (x - 1/2) * (x - 1/2))))$

using *d-gt-0* **apply** (*auto simp add: algebra-simps*)

by (*metis (no-types, opaque-lifting) add-diff-cancel-right' add-uminus-conv-diff*
cancel-comm-monoid-add-class.diff-cancel less-eq-real-def linorder-not-le mult.left-neutral
prod.simps(1))

have *v*: $(1::int, \lambda y. (d/2), 0) \neq (1, \lambda x. ((x - 1/2) * d, -(d * sqrt(1/4 - (x - 1/2) * (x - 1/2))))$

using *d-gt-0* **apply** (*auto simp add: algebra-simps*)

by (*metis (no-types, opaque-lifting) diff-0 equal-neg-zero mult-zero-left nonzero-mult-div-cancel-left*
order-less-irrefl prod.sel(1) times-divide-eq-right zero-neq-numeral)

show $card\ \{(-1::int, \lambda y. -(d/2), 0), (1, \lambda y. (d/2), 0), (1, \lambda x. ((x - 1/2) * d, -(d * sqrt(1/4 - (x - 1/2) * (x - 1/2))))),$

$(-1, \lambda x. ((x - 1/2) * d, d * sqrt(1/4 - (x - 1/2) * (x - 1/2))))\} = 4$

using *iv v* **by** *auto*

qed

lemma *rot-circle-cube-valid-two-cube*:

shows *valid-two-cube rot-circle-cube*

using *valid-cube-valid-swap circle-cube-valid-two-cube d-gt-0 rot-circle-cube-def*

by *force*

definition *circle-arc-0* **where** $circle\ arc\ 0 = (1, \lambda t::real. (0,0))$

lemma *circle-top-bot-edges-neq' [simp]*:

shows $circle\ top\ edge \neq circle\ bot\ edge$

by (simp add: circle-top-edge-def circle-bot-edge-def)

lemma *rot-circle-top-left-edges-neq* [simp]: *rot-circle-top-edge* \neq *rot-circle-left-edge*
apply (simp add: rot-circle-left-edge-def rot-circle-top-edge-def x-coord-def)
by (metis (mono-tags, opaque-lifting) cancel-comm-monoid-add-class.diff-cancel d-gt-0 divide-eq-0-iff mult-zero-left order-less-irrefl prod.sel(2) zero-neq-numeral)

lemma *rot-circle-bot-left-edges-neq* [simp]: *rot-circle-bot-edge* \neq *rot-circle-left-edge*
by (simp add: rot-circle-left-edge-def rot-circle-bot-edge-def x-coord-def)

lemma *rot-circle-top-right-edges-neq* [simp]: *rot-circle-top-edge* \neq *rot-circle-right-edge*
by (simp add: rot-circle-right-edge-def rot-circle-top-edge-def x-coord-def)

lemma *rot-circle-bot-right-edges-neq* [simp]: *rot-circle-bot-edge* \neq *rot-circle-right-edge*
apply (simp add: rot-circle-right-edge-def rot-circle-bot-edge-def x-coord-def)
by (metis (mono-tags, opaque-lifting) cancel-comm-monoid-add-class.diff-cancel d-gt-0 divide-eq-0-iff mult-zero-left neg-0-equal-iff-equal order-less-irrefl prod.sel(2) zero-neq-numeral)

lemma *rot-circle-right-top-edges-neq'* [simp]: *rot-circle-right-edge* \neq *rot-circle-left-edge*
by (simp add: rot-circle-left-edge-def rot-circle-right-edge-def)

lemma *rot-circle-left-bot-edges-neq* [simp]: *rot-circle-left-edge* \neq *rot-circle-top-edge*
apply (simp add: rot-circle-top-edge-def rot-circle-left-edge-def)
by (metis (no-types, opaque-lifting) cancel-comm-monoid-add-class.diff-cancel d-gt-0 mult.commute mult-zero-right nonzero-mult-div-cancel-left order-less-irrefl prod.sel(2) times-divide-eq-right x-coord-def zero-neq-numeral)

lemma *circle-right-top-edges-neq* [simp]: *circle-right-edge* \neq *circle-top-edge*
proof –
have *circle-right-edge* = (1, ($\lambda r::\text{real}.$ (d / 2, 0::real)))
using *circle.circle-right-edge-def circle-axioms* **by** blast
then show ?thesis
using *circle.circle-top-edge-def circle-axioms* **by** auto
qed

lemma *circle-left-bot-edges-neq* [simp]: *circle-left-edge* \neq *circle-bot-edge*
proof –
have *circle-bot-edge* = (1, $\lambda r.$ (x-coord r * d, - d * circle-y (x-coord r)))
using *circle.circle-bot-edge-def circle-axioms* **by** blast
then show ?thesis
by (simp add: circle-left-edge-def)
qed

lemma *circle-left-top-edges-neq* [simp]: *circle-left-edge* \neq *circle-top-edge*
proof –
have $\exists r.$ ((r - 1 / 2) * d, d * sqrt (1 / 4 - (r - 1 / 2) * (r - 1 / 2))) \neq
(- (d / 2), 0)
by (metis circ-left-edge-neq-top)

then have $(\exists r. d * r \neq - (d / 2)) \vee (\exists r \text{ ra. } (x\text{-coord } (x\text{-coord-inv } r) * d, d * \text{circle-y } (x\text{-coord } (x\text{-coord-inv } r))) = (x\text{-coord } \text{ra} * d, d * \text{circle-y } (x\text{-coord } \text{ra})) \wedge d * \text{circle-y } r \neq 0)$
by *(metis mult.commute)*
then have $(\lambda r. (x\text{-coord } r * d, d * \text{circle-y } (x\text{-coord } r))) \neq (\lambda r. (- (d / 2), 0))$
by *(metis mult.commute prod.simps(1) x-coord-inv-2)*
then show *?thesis*
by *(simp add: circle-left-edge-def circle-top-edge-def)*
qed

lemma *circle-right-bot-edges-neq* [*simp*]: *circle-right-edge* \neq *circle-bot-edge*
by *(smt Pair-inject circle-bot-edge-def d-gt-0 circle.circle-right-edge-def circle-axioms mult-le-cancel-right-pos x-coord-def)*

definition *circle-polar* **where**

circle-polar $t = ((d/2) * \cos (2 * \pi * t), (d/2) * \sin (2 * \pi * t))$

lemma *circle-polar-smooth*: *(circle-polar)* *C1-differentiable-on* $\{0..1\}$

proof –

have *inj* $((*) (2 * \pi))$
by *(auto simp: inj-on-def)*
then have $*$: $\bigwedge x. \text{finite } (\{0..1\} \cap (*) (2 * \pi) - \{x\})$
by *(simp add: finite-vimageI)*
have $(\lambda t. ((d/2) * \cos (2 * \pi * t), (d/2) * \sin (2 * \pi * t)))$ *C1-differentiable-on* $\{0..1\}$
by *(rule * derivative-intros)+*
then show *?thesis*
apply *(rule eq-smooth-gen)*
by *(simp add: circle-polar-def)*
qed

abbreviation *custom-arccos* $\equiv (\lambda x. (\text{if } -1 \leq x \wedge x \leq 1 \text{ then arccos } x \text{ else } (\text{if } x < -1 \text{ then } -x + \pi \text{ else } 1 - x)))$

lemma *cont-custom-arccos*:

assumes $S \subseteq \{-1..1\}$

shows *continuous-on* S *custom-arccos*

proof –

have *continuous-on* $(\{-1..1\} \cup \{\})$ *custom-arccos*

by *(auto intro!: continuous-on-cases continuous-intros)*

with *assms* **show** *?thesis*

using *continuous-on-subset* **by** *auto*

qed

lemma *custom-arccos-has-deriv*:

assumes $-1 < x < 1$

shows *DERIV* *custom-arccos* $x := \text{inverse } (- \text{sqrt } (1 - x^2))$

proof –

have $x1: |x|^2 < 1^2$

```

  by (simp add: abs-less-iff abs-square-less-1 assms)
show ?thesis
  apply (rule DERIV-inverse-function [where f=cos and a=-1 and b=1])
    apply (rule x1 derivative-eq-intros | simp add: sin-arccos)+
  using assms x1 cont-custom-arccos [of {-1<.. $1$ }]
    apply (auto simp: continuous-on-eq-continuous-at greaterThanLessThan-subseteq-atLeastAtMost-iff)
  done
qed

```

declare

```

  custom-arccos-has-deriv[THEN DERIV-chain2, derivative-intros]
  custom-arccos-has-deriv[THEN DERIV-chain2, unfolded has-field-derivative-def,
  derivative-intros]

```

lemma *circle-boundary-reparams:*

```

  shows rot-circ-left-edge-reparam-polar-circ-split:
    reparam (rec-join [(rot-circle-left-edge)]) (rec-join [(subcube (1/4) (1/2) (1,
  circle-polar)), (subcube (1/2) (3/4) (1, circle-polar))])
    (is ?P1)
  and circ-top-edge-reparam-polar-circ-split:
    reparam (rec-join [(circle-top-edge)]) (rec-join [(subcube 0 (1/4) (1, circle-polar)),
  (subcube (1/4) (1/2) (1, circle-polar))])
    (is ?P2)
  and circ-bot-edge-reparam-polar-circ-split:
    reparam (rec-join [(circle-bot-edge)]) (rec-join [(subcube (1/2) (3/4) (1, cir-
  cle-polar)), (subcube (3/4) 1 (1, circle-polar))])
    (is ?P3)
  and rot-circ-right-edge-reparam-polar-circ-split:
    reparam (rec-join [(rot-circle-right-edge)]) (rec-join [(subcube (3/4) 1 (1, cir-
  cle-polar)), (subcube 0 (1/4) (1, circle-polar))])
    (is ?P4)

```

proof –

```

  let ?φ = ((* (1/pi) ∘ custom-arccos ∘ (λt. 2 * x-coord (1 - t)))
  let ?LHS1 = (λx. (- (d * sqrt (1/4 - x-coord (1 - x) * x-coord (1 - x))),
  x-coord (1 - x) * d))
  let ?RHS1 = ((λx. if x * 2 ≤ 1 then (d * cos (2 * pi * (2 * x/4 + 1/4))/2, d
  * sin (2 * pi * (2 * x/4 + 1/4))/2)
    else (d * cos (2 * pi * ((2 * x - 1)/4 + 1/2))/2, d * sin (2
  * pi * ((2 * x - 1)/4 + 1/2))/2)) ∘ ?φ)
  let ?LHS2 = λx. ((x-coord (1 - x) * d, d * sqrt (1/4 - x-coord (1 - x) *
  x-coord (1 - x))))
  let ?RHS2 = ((λx. if x * 2 ≤ 1 then (d * cos (2 * x * pi/2)/2, d * sin (2 * x
  * pi/2)/2) else (d * cos (2 * pi * ((2 * x - 1)/4 + 1/4))/2, d * sin (2 * pi *
  ((2 * x - 1)/4 + 1/4))/2)) ∘ ?φ)
  let ?LHS3 = λx. (x-coord x * d, - (d * sqrt (1/4 - x-coord x * x-coord x)))
  let ?RHS3 = (λx. if x * 2 ≤ 1 then (d * cos (2 * pi * (2 * x/4 + 1/2))/2, d
  * sin (2 * pi * (2 * x/4 + 1/2))/2)
    else (d * cos (2 * pi * ((2 * x - 1)/4 + 3/4))/2, d * sin (2 * pi
  * ((2 * x - 1)/4 + 3/4))/2))

```

```

let ?LHS4 = λx. (d * sqrt (1/4 - x-coord x * x-coord x), x-coord x * d)
let ?RHS4 = (λx. if x * 2 ≤ 1 then (d * cos (2 * pi * (2 * x/4 + 3/4))/2, d
* sin (2 * pi * (2 * x/4 + 3/4))/2)
  else (d * cos ((2 * x - 1) * pi/2)/2, d * sin ((2 * x - 1) * pi/2)/2))
have phi-diff: ?φ piecewise-C1-differentiable-on {0..1}
  unfolding piecewise-C1-differentiable-on-def
proof
  show continuous-on {0..1} ?φ
    unfolding x-coord-def
    by (intro continuous-on-compose cont-custom-arccos continuous-intros) auto
  have ?φ C1-differentiable-on {0..1} - {0,1}
    unfolding x-coord-def piecewise-C1-differentiable-on-def C1-differentiable-on-def
  valid-path-def
    by (simp | rule has-vector-derivative-pair-within DERIV-image-chain deriva-
  tive-eq-intros continuous-intros exI ballI conjI
    | force simp add: field-simps | sos)+
  then show ∃ s. finite s ∧ ?φ C1-differentiable-on {0..1} - s
    by force
qed
have inj: inj ?φ
  apply (intro comp-inj-on inj-on-cases inj-on-arccos)
  apply (auto simp add: inj-on-def x-coord-def)
  using pi-ge-zero apply auto[1]
  apply (smt arccos)
  by (smt arccos-lbound)
then have fin: ∧x. [0 ≤ x; x ≤ 1] ⇒ finite (?φ - ' {x})
  by (simp add: finite-vimageI)
have ?φ ' {0..1} = {0..1}
proof
  show ?φ ' {0..1} ⊆ {0..1}
    by (auto simp add: x-coord-def divide-simps arccos-lbound arccos-bounded)
  have arccos (1 - 2 * ((1 - cos (x * pi))/2)) = x * pi if 0 ≤ x x ≤ 1 for x
    using that by (simp add: field-simps arccos-cos)
  then show {0..1} ⊆ ?φ ' {0..1}
    apply (auto simp: x-coord-def o-def image-def)
    by (rule-tac x=(1 - cos (x * pi))/2 in bexI) auto
qed
then have bij-phi: bij-betw ?φ {0..1} {0..1}
  unfolding bij-betw-def using inj inj-on-subset by blast
have phi01: ?φ - ' {0..1} ⊆ {0..1}
  by (auto simp add: subset-iff x-coord-def divide-simps)
{
  fix x::real assume x: 0 ≤ x x ≤ 1
  then have i: - 1 ≤ (2 * x - 1) (2 * x - 1) ≤ 1 by auto
  have ii: cos (pi / 2 + arccos (1 - x * 2)) = -sin (arccos (1 - x * 2))
    using minus-sin-cos-eq[symmetric, where ?x = arccos (1 - x * 2)]
    by (auto simp add: add commute)
  have 2 * sqrt (x - x * x) = sqrt (4*x - 4*x * x)
    by (metis mult.assoc real-sqrt-four real-sqrt-mult right-diff-distrib)
}

```

```

also have ... = sqrt (1 - (2 * x - 1) * (2 * x - 1))
  by (simp add: algebra-simps)
finally have iii: 2 * sqrt (x - x * x) = cos (arcsin (2 * x - 1)) ^ 2 * sqrt
(x - x * x) = sin (arccos (1 - 2 * x))
  using arccos-minus[OF i] unfolding minus-diff-eq sin-pi-minus
  by (simp add: cos-arcsin i power2-eq-square sin-arccos)
then have 1: ?LHS1 x = ?RHS1 x
  using d-gt-0 i x apply (simp add: x-coord-def field-simps)
  apply (auto simp add: diff-divide-distrib add-divide-distrib right-diff-distrib
mult.commute ii)
  using cos-sin-eq[where ?x = - arccos (1 - x * 2)]
  by (simp add: cos-sin-eq[where ?x = - arccos (1 - x * 2)] right-diff-distrib)
  have 2: ?LHS2 x = ?RHS2 x
    using x d-gt-0 iii by (auto simp add: x-coord-def diff-divide-distrib alge-
bra-simps)
  have a: cos (pi / 2 - arccos (x * 2 - 1)) = cos (pi / 2 - arccos (1 - x * 2))
    by (smt arccos-minus arccos-cos2 arccos-lbound cos-arccos cos-ge-minus-one
cos-le-one i(1) i(2) pi-def pi-half)
  have b: cos (arccos (1 - x * 2) + pi * 3 / 2) = cos ((pi / 2 - arccos (x * 2
- 1)) + 2 * pi)
    using arccos-minus[OF i] by (auto simp add: mult.commute add.commute)
  then have c: ... = cos (pi / 2 - arccos (x * 2 - 1)) using cos-periodic by
blast
  have cos (- pi - arccos (1 - x * 2)) = cos (- (pi + arccos (1 - x * 2)))
    by (auto simp add: minus-add[where b = pi and a = arccos (1 - x * 2),
symmetric])
  moreover have ... = cos ((pi + arccos (1 - x * 2)))
    using cos-minus by blast
  moreover have ... = cos (2*pi - arccos (x * 2 - 1))
    using arccos-minus[OF i] by (auto simp add: mult.commute add.commute)
  moreover have ... = cos (arccos (x * 2 - 1))
    using cos-2pi-minus by auto
  ultimately have d: cos (- pi - arccos (1 - x * 2)) = (x * 2 - 1)
    using cos-arccos[OF i] mult.commute by metis
  have cosm:  $\bigwedge x. \cos (x - \pi * 2) = \cos x$ 
    by (metis cos-periodic eq-diff-eq' mult.commute)
  have 34: ?LHS3 x = (?RHS3  $\circ$  ? $\varphi$ ) x ?LHS4 x = (?RHS4  $\circ$  ? $\varphi$ ) x
    using d-gt-0 x a b c iii cos-periodic [of pi / 2 - arccos (x * 2 - 1)]
  apply (auto simp add: x-coord-def algebra-simps diff-divide-distrib power2-eq-square)
  apply (auto simp add: sin-cos-eq cosm)
  using d apply (auto simp add: right-diff-distrib)
  by (smt cos-minus)
note 1 2 34
} note * = this
show ?P1 ?P2 ?P3 ?P4
  apply (auto simp add: subcube-def circle-bot-edge-def circle-top-edge-def
circle-polar-def reversepath-def
  subpath-def joinpaths-def circle-y-def rot-circle-left-edge-def rot-circle-right-edge-def)
  unfolding reparam-def

```

by (rule ballI exI conjI impI phi-diff bij-phi phi01 fin * | force simp add:
x-coord-def)+
qed

definition *circle-cube-boundary-to-polarcircle* **where**

circle-cube-boundary-to-polarcircle $\gamma \equiv$
if ($\gamma = (\text{circle-top-edge})$) then
 {subcube 0 (1/4) (1, circle-polar), subcube (1/4) (1/2) (1, circle-polar)}
else if ($\gamma = (\text{circle-bot-edge})$) then
 {subcube (1/2) (3/4) (1, circle-polar), subcube (3/4) 1 (1, circle-polar)}
else {}

definition *rot-circle-cube-boundary-to-polarcircle* **where**

rot-circle-cube-boundary-to-polarcircle $\gamma \equiv$
if ($\gamma = (\text{rot-circle-left-edge})$) then
 {subcube (1/4) (1/2) (1, circle-polar), subcube (1/2) (3/4) (1, circle-polar)}
else if ($\gamma = (\text{rot-circle-right-edge})$) then
 {subcube (3/4) 1 (1, circle-polar), subcube 0 (1/4) (1, circle-polar)}
else {}

lemma *circle-arcs-neq*:

assumes $0 \leq k \leq 1$ $0 \leq n \leq 1$ $n < k$ $k + n < 1$
shows *subcube* $k \ m \ (1, \text{circle-polar}) \neq$ *subcube* $n \ q \ (1, \text{circle-polar})$

proof (simp add: subcube-def subpath-def circle-polar-def)

have $\cos(2 * \pi * k) \neq \cos(2 * \pi * n)$

unfolding *cos-eq*

proof safe

show *False* if $2 * \pi * k = 2 * \pi * n + 2 * m * \pi$ $m \in \mathbb{Z}$ **for** m

proof –

have $2 * \pi * (k - n) = 2 * m * \pi$

using *distrib-left* that **by** (simp add: *left-diff-distrib* *mult.commute*)

then have $a: m = (k - n)$ **by** *auto*

have $\lfloor k - n \rfloor = 0$

using *assms* **by** (simp add: *floor-eq-iff*)

then have $k - n \notin \mathbb{Z}$

using *assms* **by** (*auto simp only: frac-eq-0-iff[symmetric] frac-def*)

then show *False* using that a **by** *auto*

qed

show *False* if $2 * \pi * k = - (2 * \pi * n) + 2 * m * \pi$ $m \in \mathbb{Z}$ **for** m

proof –

have $2 * \pi * (k + n) = 2 * m * \pi$

using that **by** (*auto simp add: distrib-left*)

then have $a: m = (k + n)$ **by** *auto*

have $\lfloor k + n \rfloor = 0$

using *assms* **by** (simp add: *floor-eq-iff*)

then have $k + n \notin \mathbb{Z}$

using *Ints-def* *assms* **by** *force*

then show *False using that a by auto*
qed
qed
then have $(\lambda x. (d * \cos (2 * \pi * ((m - k) * x + k))/2, d * \sin (2 * \pi * ((m - k) * x + k))/2)) 0 \neq (\lambda x. (d * \cos (2 * \pi * ((q - n) * x + n))/2, d * \sin (2 * \pi * ((q - n) * x + n))/2)) 0$
using *d-gt-0 by auto*
then show $(\lambda x. (d * \cos (2 * \pi * ((m - k) * x + k))/2, d * \sin (2 * \pi * ((m - k) * x + k))/2)) \neq (\lambda x. (d * \cos (2 * \pi * ((q - n) * x + n))/2, d * \sin (2 * \pi * ((q - n) * x + n))/2))$
by *metis*
qed

lemma *circle-arcs-neq-2:*
assumes $0 \leq k \leq 1 \ 0 \leq n \leq 1 \ n < k \ 0 < n$ **and** *kn12: $1/2 < k + n$ and $k + n < 3/2$*
shows *subcube k m (1, circle-polar) \neq subcube n q (1, circle-polar)*
proof (*simp add: subcube-def subpath-def circle-polar-def*)
have $\sin (2 * \pi * k) \neq \sin (2 * \pi * n)$
unfolding *sin-eq*
proof *safe*
show *False if $m \in \mathbb{Z} \ 2 * \pi * k = 2 * \pi * n + 2 * m * \pi$ for m*
proof –
have $2 * \pi * (k - n) = 2 * m * \pi$
using *that by (simp add: left-diff-distrib mult.commute)*
then have *a: $m = (k - n)$ by auto*
have $\lfloor k - n \rfloor = 0$
using *assms by (simp add: floor-eq-iff)*
then have $k - n \notin \mathbb{Z}$
using *assms by (auto simp only: frac-eq-0-iff[symmetric] frac-def)*
then show *False using that a by auto*
qed
show *False if $2 * \pi * k = - (2 * \pi * n) + (2 * m + 1) * \pi \ m \in \mathbb{Z}$ for m*
proof –
have *i: $\bigwedge \pi. 0 < \pi \implies 2 * \pi * (k + n) = 2 * m * \pi + \pi \implies m = (k + n) - 1/2$*
by (*sos (((((A < 0 * A < 1) * R < 1) + ([1/2] * A = 0))) & (((A < 0 * A < 1) * R < 1) + ([~ 1/2] * A = 0))))*)
have $2 * \pi * (k + n) = 2 * m * \pi + \pi$
using *that by (simp add: algebra-simps)*
then have *a: $m = (k + n) - 1/2$ using i[OF pi-gt-zero] by fastforce*
have $\lfloor k + n - 1/2 \rfloor = 0$
using *assms by (auto simp add: floor-eq-iff)*
then have $k + n - 1/2 \notin \mathbb{Z}$
by (*metis Ints-cases add.commute add.left-neutral add-diff-cancel-left' add-diff-eq kn12 floor-of-int of-int-0 order-less-irrefl*)
then show *False using that a by auto*
qed
qed

```

then have ( $\lambda x. (d * \cos (2 * \pi * ((m - k) * x + k))/2, d * \sin (2 * \pi * ((m - k) * x + k))/2)$ )  $0 \neq$  ( $\lambda x. (d * \cos (2 * \pi * ((q - n) * x + n))/2, d * \sin (2 * \pi * ((q - n) * x + n))/2)$ )  $0$ 
using d-gt-0 by auto
then show ( $\lambda x. (d * \cos (2 * \pi * ((m - k) * x + k))/2, d * \sin (2 * \pi * ((m - k) * x + k))/2)$ )  $\neq$  ( $\lambda x. (d * \cos (2 * \pi * ((q - n) * x + n))/2, d * \sin (2 * \pi * ((q - n) * x + n))/2)$ )
by metis
qed

```

lemma *circle-cube-is-only-horizontal-div-of-rot*:

```

shows only-horizontal-division (boundary (circle-cube)) {rot-circle-cube}
unfolding only-horizontal-division-def
proof (rule exI [of - {}], simp, intro conjI ballI)
show finite (boundary circle-cube)
using circle.circle-cube-boundary-explicit circle-axioms by auto
show boundary-chain (boundary circle-cube)
by (simp add: two-cube-boundary-is-boundary)
show  $\bigwedge x. x \in \text{boundary } \text{circle-cube} \implies \text{case } x \text{ of } (k, x) \Rightarrow \text{valid-path } x$ 
using circle-cube-boundary-valid by blast
let  $?V = (\text{boundary } (\text{circle-cube}))$ 
let  $?pi = \{\text{circle-left-edge}, \text{circle-right-edge}\}$ 
let  $?pj = \{\text{rot-circle-top-edge}, \text{rot-circle-bot-edge}\}$ 
let  $?f = \text{circle-cube-boundary-to-polarcircle}$ 
let  $?one-chaini = \text{boundary } (\text{circle-cube}) - ?pi$ 
have c: common-reparam-exists  $?V$  (two-chain-vertical-boundary {rot-circle-cube})
unfolding common-reparam-exists-def
proof (intro exI conjI)
let  $?subdiv = \{(\text{subcube } 0 (1/4) (1, \text{circle-polar})),$ 
 $(\text{subcube } (1/4) (1/2) (1, \text{circle-polar})),$ 
 $(\text{subcube } (1/2) (3/4) (1, \text{circle-polar})),$ 
 $(\text{subcube } (3/4) 1 (1, \text{circle-polar}))\}$ 
show  $(\forall (k, \gamma) \in ?subdiv. \gamma \text{ C1-differentiable-on } \{0..1\})$ 
using subpath-smooth[OF circle-polar-smooth] by (auto simp add: subcube-def)
have 1: finite  $?subdiv$  by auto
show boundary-chain  $?subdiv$ 
by (simp add: boundary-chain-def subcube-def)
show chain-reparam-chain' (boundary (circle-cube) -  $?pi$ )  $?subdiv$ 
unfolding chain-reparam-chain'-def
proof (intro exI conjI impI)
show  $\bigcup (?f \text{ ` } ?one-chaini) = ?subdiv$ 
apply (simp add: circle-cube-boundary-to-polarcircle-def circle-cube-boundary-explicit)
using circle-top-bot-edges-neq' by metis
let  $?l = [\text{subcube } 0 (1/4) (1, \text{circle-polar}), \text{subcube } (1/4) (1/2) (1, \text{circle-polar})]$ 
have chain-reparam-weak-path (coeff-cube-to-path (circle-top-edge)) {subcube
 $0 (1/4) (1, \text{circle-polar}), \text{subcube } (1/4) (1/2) (1, \text{circle-polar})\}$ 
unfolding chain-reparam-weak-path-def
proof (intro exI conjI)

```

```

show valid-chain-list ?l
  by (auto simp add: subcube-def circle-top-edge-def x-coord-def circle-y-def
pathfinish-def pathstart-def
      reversepath-def subpath-def joinpaths-def)
show reparam (coeff-cube-to-path circle-top-edge) (rec-join ?l)
  using circ-top-edge-reparam-polar-circ-split by auto
show distinct ?l
  apply simp
  apply (subst neq-commute)
  apply (simp add: circle-arcs-neq)
  done
qed auto
moreover have chain-reparam-weak-path (coeff-cube-to-path (circle-bot-edge))
{subcube (1/2) (3/4) (1, circle-polar), subcube (3/4) 1 (1, circle-polar)}
  unfolding chain-reparam-weak-path-def
proof
  let ?l = [subcube (1/2) (3/4) (1, circle-polar), subcube (3/4) 1 (1,
circle-polar)]
  have a: valid-chain-list ?l
  by (auto simp add: subcube-def circle-top-edge-def x-coord-def circle-y-def
pathfinish-def pathstart-def
      reversepath-def subpath-def joinpaths-def)
  have b: reparam (rec-join [circle-bot-edge]) (rec-join ?l)
  using circ-bot-edge-reparam-polar-circ-split by auto
  have c: subcube (3/4) 1 (1, circle-polar) ≠ subcube (1/2) (3/4) (1,
circle-polar)
  apply(rule circle-arcs-neq-2) using d-gt-0(1) by auto
  show set ?l = {subcube (1/2) (3/4) (1, circle-polar), subcube (3/4) 1 (1,
circle-polar)}  $\wedge$ 
      distinct ?l  $\wedge$  reparam (coeff-cube-to-path (circle-bot-edge))
(rec-join ?l)  $\wedge$  valid-chain-list ?l  $\wedge$  ?l ≠ [] using a b c by auto
qed
ultimately
show ( $\forall$  cube  $\in$  ?one-chaini. chain-reparam-weak-path (rec-join [cube]) (?f cube))
  by (auto simp add: circle-cube-boundary-to-polarcircle-def UNION-eq cir-
cle-cube-boundary-explicit)
  show ( $\forall$  p  $\in$  ?one-chaini.  $\forall$  p'  $\in$  ?one-chaini. p  $\neq$  p'  $\longrightarrow$  ?f p  $\cap$  ?f p' = {})
  using circle-arcs-neq circle-arcs-neq-2
  apply (auto simp add: circle-cube-boundary-to-polarcircle-def UNION-eq
circle-cube-boundary-explicit neq-commute d-gt-0)
  using circle-top-bot-edges-neq' d-gt-0 apply auto[1]
  apply (smt atLeastAtMost-iff divide-less-eq-1-pos zero-less-divide-1-iff)
  apply (smt atLeastAtMost-iff divide-less-eq-1-pos zero-less-divide-iff)
  apply (smt atLeastAtMost-iff divide-cancel-left divide-less-eq-1-pos field-sum-of-halves
zero-less-divide-1-iff)
  done
show ( $\forall$  x  $\in$  ?one-chaini. finite (?f x))
by (auto simp add: circle-cube-boundary-to-polarcircle-def circle-cube-boundary-explicit)
qed

```

```

show ( $\forall (k, \gamma) \in ?pi. \text{point-path } \gamma$ )
using d-gt-0 by (auto simp add: point-path-def circle-left-edge-def circle-right-edge-def)
let  $?f = \text{rot-circle-cube-boundary-to-polarcircle}$ 
let  $?one-chain2.0 = \text{two-chain-vertical-boundary } \{\text{rot-circle-cube}\} - ?pj$ 
show  $\text{chain-reparam-chain}' (\text{two-chain-vertical-boundary } \{\text{rot-circle-cube}\} -$ 
 $?pj) ?subdiv$ 
unfolding chain-reparam-chain'-def
proof (intro exI conjI)
have  $rw: ?one-chain2.0 = \{\text{rot-circle-left-edge}, \text{rot-circle-right-edge}\}$ 
by (auto simp add: rot-circle-cube-vertical-boundary-explicit two-chain-vertical-boundary-def)
show  $\bigcup (?f' ?one-chain2.0) = ?subdiv$ 
using rot-circle-right-top-edges-neq'
by (auto simp add: rot-circle-cube-boundary-to-polarcircle-def rw)
show ( $\forall \text{cube} \in ?one-chain2.0. \text{chain-reparam-weak-path } (\text{rec-join } [\text{cube}]) (?f$ 
 $\text{cube}))$ 
proof (clarsimp simp add: rot-circle-cube-boundary-to-polarcircle-def rw, intro
 $\text{conjI}$ )
let  $?l = [\text{subcube } (1/4) (1/2) (1, \text{circle-polar}), \text{subcube } (1/2) (3/4) (1,$ 
 $\text{circle-polar})]$ 
show  $\text{chain-reparam-weak-path } (\text{coeff-cube-to-path } (\text{rot-circle-left-edge}))$ 
 $\{\text{subcube } (1/4) (1/2) (1, \text{circle-polar}), \text{subcube } (1/2) (3/4) (1, \text{circle-polar})\}$ 
unfolding chain-reparam-weak-path-def
proof (intro exI conjI)
show valid-chain-list ?l
by (auto simp add: subcube-def pathfinish-def pathstart-def reversepath-def
 $\text{subpath-def joinpaths-def}$ )
show  $\text{reparam } (\text{coeff-cube-to-path } \text{rot-circle-left-edge}) (\text{rec-join } ?l)$ 
using rot-circ-left-edge-reparam-polar-circ-split by auto
show distinct ?l
apply simp
apply (subst neq-commute)
apply (simp add: circle-arcs-neq)
done
qed auto
show  $\text{chain-reparam-weak-path } (\text{coeff-cube-to-path } (\text{rot-circle-right-edge}))$ 
 $\{\text{subcube } (3/4) 1 (1, \text{circle-polar}), \text{subcube } 0 (1/4) (1, \text{circle-polar})\}$ 
unfolding chain-reparam-weak-path-def
proof (intro exI conjI)
let  $?l = [\text{subcube } (3/4) 1 (1, \text{circle-polar}), \text{subcube } 0 (1/4) (1, \text{circle-polar})]$ 
show valid-chain-list ?l
by (auto simp add: circle-polar-def subcube-def pathfinish-def pathstart-def
 $\text{reversepath-def subpath-def joinpaths-def}$ )
show  $\text{reparam } (\text{coeff-cube-to-path } \text{rot-circle-right-edge}) (\text{rec-join } ?l)$ 
using rot-circ-right-edge-reparam-polar-circ-split by auto
show distinct ?l
by (simp add: circle-arcs-neq)
qed auto
qed
show ( $\forall p \in ?one-chain2.0. \forall p' \in ?one-chain2.0. p \neq p' \longrightarrow ?f p \cap ?f p' = \{\}$ )

```

```

using circle-arcs-neq circle-arcs-neq-2
apply (auto simp add: rot-circle-cube-boundary-to-polarcircle-def neq-commute)
apply (metis add.right-neutral divide-less-eq-1-pos dual-order.order-iff-strict
num.distinct(1) one-less-numeral-iff prod.sel(1) prod.sel(2) semiring-norm(68) sub-
cube-def zero-less-divide-1-iff zero-less-numeral)
apply (smt field-sum-of-halves)
done
show ( $\forall x \in ?one-chain2.0. finite (?f x)$ )
by (auto simp add: rot-circle-cube-boundary-to-polarcircle-def UNION-eq rw)
qed
show ( $\forall (k, \gamma) \in ?pj. point-path \gamma$ )
using d-gt-0 by (auto simp add: point-path-def rot-circle-top-edge-def rot-circle-bot-edge-def)
qed
then show common-sudiv-exists (two-chain-vertical-boundary {rot-circle-cube})
(boundary circle-cube)  $\vee$ 
common-reparam-exists (boundary circle-cube) (two-chain-vertical-boundary
{rot-circle-cube})
by blast
qed

```

lemma *GreenThm-cirlce*:

```

assumes  $\forall twoC \in \{circle-cube\}. analytically-valid (cubeImage twoC) (\lambda x. F x \cdot$ 
i) j
 $\forall twoC \in \{rot-circle-cube\}. analytically-valid (cubeImage twoC) (\lambda x. F x \cdot j) i$ 
shows integral (cubeImage (circle-cube)) ( $\lambda x. partial-vector-derivative (\lambda x. F x \cdot$ 
j) i x - partial-vector-derivative ( $\lambda x. F x \cdot i) j x$ ) =
one-chain-line-integral F {i, j} (boundary (circle-cube))
proof(rule green-typeI-typeII-chain.GreenThm-typeI-typeII-divisible-region-finite-holes[of
(cubeImage (circle-cube)) i j F {circle-cube} {rot-circle-cube},
OF - - - circle-cube-is-only-horizontal-div-of-rot -], auto)
show  $\wedge a b. (a, b) \in boundary circle-cube \implies valid-path b$  using circle-cube-boundary-valid
by auto
show green-typeI-typeII-chain (cubeImage circle-cube) i j F {circle-cube} {rot-circle-cube}
using assms
proof(auto simp add: green-typeI-typeII-chain-def green-typeI-chain-def green-typeII-chain-def
green-typeI-chain-axioms-def green-typeII-chain-axioms-def
intro!: circle-cube-is-type-I rot-circle-cube-is-type-II d-gt-0 R2-axioms)
show gen-division (cubeImage circle-cube) {cubeImage circle-cube} by (simp
add: gen-division-def)
show gen-division (cubeImage (circle-cube)) ({cubeImage rot-circle-cube})
using rot-circle-div-circle d-gt-0 by auto
show valid-two-chain {rot-circle-cube} valid-two-chain {circle-cube}
apply (auto simp add: valid-two-chain-def)
using rot-circle-cube-valid-two-cube circle-cube-valid-two-cube assms(1) by
auto
qed
show only-vertical-division (boundary (circle-cube)) {circle-cube}
using twoChainVertDiv-of-itself[of {circle-cube}]
apply(simp add: two-chain-boundary-def)

```

```

    using circle-cube-boundary-valid
    by auto
qed
end
end

```

3 The Diamond Example

```

theory DiamExample
  imports Green SymmetricR2Shapes
begin

```

```

lemma abs-if':
  fixes a :: 'a :: {abs-if,ordered-ab-group-add}
  shows |a| = (if a ≤ 0 then - a else a)
  by (simp add: abs-if dual-order.order-iff-strict)

```

```

locale diamond = R2 +
  fixes d::real
  assumes d-gt-0: 0 < d
begin

```

```

definition diamond-y-gen :: real ⇒ real where
  diamond-y-gen ≡ λt. 1/2 - |t|

```

```

definition diamond-cube-gen:: ((real * real) ⇒ (real * real)) where
  diamond-cube-gen ≡ (λ(x,y). (d * x-coord x, (2 * y - 1) * (d * diamond-y-gen
(x-coord x))))

```

```

lemma diamond-y-gen-valid:
  assumes a ≤ 0 0 ≤ b
  shows diamond-y-gen piecewise-C1-differentiable-on {a..b}
  unfolding piecewise-C1-differentiable-on-def diamond-y-gen-def
proof (intro conjI exI)
  show continuous-on {a..b} (λt. 1 / 2 - |t|)
    by (intro continuous-intros)
  show finite{0}
    by simp
  show (λt. 1 / 2 - |t|) C1-differentiable-on {a..b} - {0}
    by (intro derivative-intros) auto
qed

```

```

lemma diamond-cube-gen-boundary-valid:
  assumes (k,γ)∈boundary (diamond-cube-gen)
  shows valid-path γ
  using assms
proof (auto simp add: valid-path-def boundary-def horizontal-boundary-def verti-
cal-boundary-def diamond-cube-gen-def)
  have rw2: (λx. diamond-y-gen (x-coord x)) = diamond-y-gen o x-coord by auto

```

note $[derivative-intros] = C1\text{-differentiable-on-pair pair-prod-smooth-pw-smooth}$
scale-piecewise-C1-differentiable-on piecewise-C1-differentiable-neg piecewise-C1-differentiable-compose
diamond-y-gen-valid
show $(\lambda x. (d * x\text{-coord } x, - (d * \text{diamond-y-gen } (x\text{-coord } x))))$ *piecewise-C1-differentiable-on*
 $\{0..1\}$
apply *(auto intro!: derivative-intros)+*
apply *(auto simp add: x-coord-smooth rw2)*
by *(auto intro!: derivative-intros simp add: x-coord-img x-coord-back-img C1-differentiable-imp-piecewise[OF*
x-coord-smooth])+
show $(\lambda x. (d * x\text{-coord } x, d * \text{diamond-y-gen } (x\text{-coord } x)))$ *piecewise-C1-differentiable-on*
 $\{0..1\}$
apply *(auto intro!: derivative-intros)+*
apply *(auto simp add: x-coord-smooth rw2)*
by *(auto intro!: derivative-intros simp add: x-coord-img x-coord-back-img C1-differentiable-imp-piecewise[OF*
x-coord-smooth])+
qed

definition *diamond-x* **where**
 $diamond-x \equiv \lambda t. (t - 1/2) * d$

definition *diamond-y* **where**
 $diamond-y \equiv \lambda t. d/2 - |t|$

definition *diamond-cube* **where**
 $diamond-cube = (\lambda(x,y). (diamond-x\ x, (2 * y - 1) * (diamond-y\ (diamond-x\ x))))$

definition *rot-diamond-cube* **where**
 $rot-diamond-cube = prod.swap\ o\ (diamond-cube)\ o\ prod.swap$

lemma *diamond-eq-characterisations*:
shows $diamond-cube\ (x,y) = diamond-cube-gen\ (x,y)$
by *(auto simp add: diamond-cube-def diamond-cube-gen-def diamond-x-def x-coord-def*
diamond-y-def diamond-y-gen-def d-gt-0 field-simps mult-le-0-iff abs-if split: if-split-asm)

lemma *diamond-eq-characterisations-fun*: $diamond-cube = diamond-cube-gen$
using *diamond-eq-characterisations* **by** *auto*

lemma *diamond-y-valid*:
shows $diamond-y$ *piecewise-C1-differentiable-on* $\{-d/2..d/2\}$ **(is ?P)**
 $(\lambda x. diamond-y\ x)$ *piecewise-C1-differentiable-on* $\{-d/2..d/2\}$ **(is ?Q)**

proof –
have $f0: finite\ \{0\}$
by *simp*
show $?P\ ?Q$
unfolding *piecewise-C1-differentiable-on-def diamond-y-def*
by *(fastforce intro!: f0 continuous-intros derivative-intros)+*
qed

```

lemma diamond-cube-boundary-valid:
  assumes  $(k, \gamma) \in \text{boundary } (\text{diamond-cube})$ 
  shows valid-path  $\gamma$ 
  using diamond-cube-gen-boundary-valid assms d-gt-0
  by (simp add: diamond-eq-characterisations-fun)

lemma diamond-cube-is-type-I:
  shows typeI-twoCube (diamond-cube)
  unfolding typeI-twoCube-def
proof (intro exI conjI ballI)
  show  $-d/2 < d/2$ 
    using d-gt-0 by auto
  show  $\bigwedge x. x \in \{-d/2..d/2\} \implies -\text{diamond-y } x \leq \text{diamond-y } x$ 
    using diamond-y-def by auto
  have f0: finite  $\{0\}$ 
    by simp
  show diamond-y piecewise-C1-differentiable-on  $\{-d/2..d/2\}$ 
     $(\lambda x. -\text{diamond-y } x)$  piecewise-C1-differentiable-on  $\{-d/2..d/2\}$ 
    unfolding diamond-y-def piecewise-C1-differentiable-on-def
    by (rule conjI exI f0 continuous-intros derivative-intros | force)+
  show diamond-cube =
     $(\lambda(x, y). ((1 - x) * (-d/2) + x * (d/2),$ 
       $(1 - y) * -\text{diamond-y } ((1 - x) * (-d/2) + x * (d/2)) +$ 
       $y * \text{diamond-y } ((1 - x) * (-d/2) + x * (d/2))))$ 
    by (auto simp: diamond-cube-def diamond-x-def diamond-y-def divide-simps
algebra-simps)
qed

lemma diamond-cube-valid-two-cube:
  shows valid-two-cube (diamond-cube)
  apply (auto simp add: valid-two-cube-def boundary-def horizontal-boundary-def
vertical-boundary-def diamond-cube-def
diamond-x-def card-insert-if)
  apply (metis (no-types, opaque-lifting) cancel-comm-monoid-add-class.diff-cancel
d-gt-0 mult.commute mult-2 mult-zero-right order-less-irrefl prod.inject field-sum-of-halves)
  apply (metis (no-types, opaque-lifting) add-diff-cancel-right' d-gt-0 mult-cancel-left
mult-zero-right order-less-irrefl prod.inject)
  done

lemma rot-diamond-cube-boundary-valid:
  assumes  $(k, \gamma) \in \text{boundary } (\text{rot-diamond-cube})$ 
  shows valid-path  $\gamma$ 
  using d-gt-0 swap-valid-boundaries diamond-cube-boundary-valid
  using assms diamond-cube-is-type-I rot-diamond-cube-def by fastforce

lemma rot-diamond-cube-is-type-II:
  shows typeII-twoCube (rot-diamond-cube)
  using d-gt-0 swap-typeI-is-typeII diamond-cube-is-type-I
  by (auto simp add: rot-diamond-cube-def)

```


lemma *rot-diamond-cube-valid-two-cube*: *valid-two-cube* (*rot-diamond-cube*)
using *valid-cube-valid-swap diamond-cube-valid-two-cube d-gt-0 rot-diamond-cube-def*
by *force*

definition *diamond-top-edges* **where**

$diamond\text{-top-edges} = (- 1::int, \lambda x. (diamond\text{-}x\ x, diamond\text{-}y\ (diamond\text{-}x\ x)))$

definition *diamond-bot-edges* **where**

$diamond\text{-bot-edges} = (1::int, \lambda x. (diamond\text{-}x\ x, -\ diamond\text{-}y\ (diamond\text{-}x\ x)))$

lemma *diamond-cube-boundary-explicit*:

$boundary\ (diamond\text{-}cube) =$

$\{diamond\text{-}top\text{-}edges,$

$diamond\text{-}bot\text{-}edges,$

$(- 1::int, \lambda y. (diamond\text{-}x\ 0, (2 * y - 1) * diamond\text{-}y\ (diamond\text{-}x\ 0))),$

$(1::int, \lambda y. (diamond\text{-}x\ 1, (2 * y - 1) * diamond\text{-}y\ (diamond\text{-}x\ 1)))\}$

by (*auto simp only: diamond-top-edges-def diamond-bot-edges-def valid-two-cube-def boundary-def horizontal-boundary-def vertical-boundary-def diamond-cube-def Un-iff algebra-simps*)

definition *diamond-top-left-edge* **where**

$diamond\text{-top-left-edge} = (- 1::int, (\lambda x. (diamond\text{-}x\ (1/2 * x), (diamond\text{-}x\ (1/2 * x)) + d/2)))$

definition *diamond-top-right-edge* **where**

$diamond\text{-top-right-edge} = (- 1::int, (\lambda x. (diamond\text{-}x\ (1/2 * x + 1/2), -(diamond\text{-}x\ (1/2 * x + 1/2)) + d/2)))$

definition *diamond-bot-left-edge* **where**

$diamond\text{-bot-left-edge} = (1::int, (\lambda x. (diamond\text{-}x\ (1/2 * x), -(diamond\text{-}x\ (1/2 * x)) + d/2)))$

definition *diamond-bot-right-edge* **where**

$diamond\text{-bot-right-edge} = (1::int, (\lambda x. (diamond\text{-}x\ (1/2 * x + 1/2), -(-(diamond\text{-}x\ (1/2 * x + 1/2)) + d/2)))$

lemma *diamond-edges-are-valid*:

$valid\text{-}path\ (snd\ (diamond\text{-}top\text{-}left\text{-}edge))$

$valid\text{-}path\ (snd\ (diamond\text{-}top\text{-}right\text{-}edge))$

$valid\text{-}path\ (snd\ (diamond\text{-}bot\text{-}left\text{-}edge))$

$valid\text{-}path\ (snd\ (diamond\text{-}bot\text{-}right\text{-}edge))$

by (*auto simp add: valid-path-def diamond-top-left-edge-def diamond-bot-right-edge-def diamond-bot-left-edge-def diamond-top-right-edge-def diamond-x-def*)

definition *diamond-cube-boundary-to-subdiv* **where**

$diamond\text{-cube-boundary-to-subdiv}\ (\gamma::(int \times (real \Rightarrow real \times real))) \equiv$

if ($\gamma = diamond\text{-}top\text{-}edges$) *then*

$\{diamond\text{-}top\text{-}left\text{-}edge, diamond\text{-}top\text{-}right\text{-}edge\}$

else if (gamma = diamond-bot-edges) then
 {diamond-bot-left-edge, diamond-bot-right-edge}
 else {}

lemma rot-diam-edge-1:

(1::int, λx::real. ((x::real) * (2 * diamond-y (diamond-x 0)) - 1 * diamond-y (diamond-x 0), diamond-x 0) =
 (1, λx. (x * (2 * diamond-y (diamond-x 0)) - (diamond-y (diamond-x 0)),
 diamond-x 0))
 by (auto simp add: algebra-simps)

definition diamond-left-edges where

diamond-left-edges = (- 1, λy. (- diamond-y (diamond-x y), diamond-x y))

definition diamond-right-edges where

diamond-right-edges = (1, λy. (diamond-y (diamond-x y), diamond-x y))

lemma rot-diamond-cube-boundary-explicit:

boundary (rot-diamond-cube) = {(1::int, λx::real. ((2 * x - 1) * diamond-y (diamond-x 0), diamond-x 0)),
 (- 1, λx. ((2 * x - 1) * diamond-y (diamond-x 1),
 diamond-x 1)),
 diamond-left-edges, diamond-right-edges}

proof –

have boundary (rot-diamond-cube) = { (1::int, λx::real. ((2 * x - 1) * diamond-y (diamond-x 0), diamond-x 0)),
 (- 1, λx. ((2 * x - 1) * diamond-y (diamond-x 1), diamond-x 1)),
 (- 1, λy. ((2 * 0 - 1) * diamond-y (diamond-x y), diamond-x y)),
 (1, λy. ((2 * 1 - 1) * diamond-y (diamond-x y), diamond-x y))}

by (auto simp only: valid-two-cube-def boundary-def horizontal-boundary-def vertical-boundary-def rot-diamond-cube-def diamond-cube-def o-def swap-simp Un-iff)

then show ?thesis

by (auto simp add: diamond-left-edges-def diamond-right-edges-def)

qed

lemma rot-diamond-cube-vertical-boundary-explicit:

vertical-boundary (rot-diamond-cube) = {diamond-left-edges, diamond-right-edges}

proof –

have vertical-boundary (rot-diamond-cube) = {(- 1::int, λy. ((2 * 0 - 1) * diamond-y (diamond-x y), diamond-x y)),
 (1, λy. ((2 * 1 - 1) * diamond-y (diamond-x y), diamond-x y))}

by (auto simp only: valid-two-cube-def boundary-def horizontal-boundary-def vertical-boundary-def rot-diamond-cube-def diamond-cube-def o-def swap-simp Un-iff)

then show ?thesis

by (auto simp add: diamond-left-edges-def diamond-right-edges-def)

qed

definition rot-diamond-cube-boundary-to-subdiv where

$rot\text{-}diamond\text{-}cube\text{-}boundary\text{-}to\text{-}subdiv (gamma::(int \times (real \Rightarrow real \times real))) \equiv$
 $if (gamma = diamond\text{-}left\text{-}edges) then \{diamond\text{-}bot\text{-}left\text{-}edge, diamond\text{-}top\text{-}left\text{-}edge\}$
 $else if (gamma = diamond\text{-}right\text{-}edges) then \{diamond\text{-}bot\text{-}right\text{-}edge, dia-$
 $mond\text{-}top\text{-}right\text{-}edge\}$
 $else \{\}$

definition *diamond-boundaries-reparam-map* **where**

diamond-boundaries-reparam-map $\equiv id$

lemma *diamond-boundaries-reparam-map-bij*:

bij (diamond-boundaries-reparam-map)

by (*auto simp add: bij-def full-SetCompr-eq[symmetric] diamond-boundaries-reparam-map-def*)

lemma *diamond-bot-edges-neq-diamond-top-edges*:

diamond-bot-edges \neq *diamond-top-edges*

by (*simp add: diamond-bot-edges-def diamond-top-edges-def*)

lemma *diamond-top-left-edge-neq-diamond-top-right-edge*:

diamond-top-left-edge \neq *diamond-top-right-edge*

apply (*simp add: diamond-top-left-edge-def diamond-top-right-edge-def diamond-x-def diamond-y-def*)

using *d-gt-0*

apply (*auto simp add: algebra-simps divide-simps*)

by (*metis (no-types, opaque-lifting) diff-zero div-0 divide-divide-eq-right order-less-irrefl prod.inject field-sum-of-halves*)

lemma *neqs1*:

shows $(\lambda x. (diamond\text{-}x\ x, diamond\text{-}y (diamond\text{-}x\ x))) \neq (\lambda x. (diamond\text{-}x\ x, - diamond\text{-}y (diamond\text{-}x\ x)))$

and $(\lambda y. (- diamond\text{-}y (diamond\text{-}x\ y), diamond\text{-}x\ y)) \neq (\lambda y. (diamond\text{-}y (diamond\text{-}x\ y), diamond\text{-}x\ y))$

and $(\lambda x. (diamond\text{-}x(x/2 + 1/2), diamond\text{-}x(x/2 + 1/2) - d/2)) \neq (\lambda x. (diamond\text{-}x(x/2), - diamond\text{-}x(x/2) - d/2))$

and $(\lambda x. (diamond\text{-}x(x/2 + 1/2), d/2 - diamond\text{-}x(x/2 + 1/2))) \neq (\lambda x. (diamond\text{-}x(x/2), diamond\text{-}x(x/2) + d/2))$

and $(\lambda x. (diamond\text{-}x(x/2), - diamond\text{-}x(x/2) - d/2)) \neq (\lambda x. (diamond\text{-}x(x/2 + 1/2), diamond\text{-}x(x/2 + 1/2) - d/2))$

and $(\lambda x. (diamond\text{-}x(x/2), diamond\text{-}x(x/2) + d/2)) \neq (\lambda x. (diamond\text{-}x(x/2 + 1/2), d/2 - diamond\text{-}x(x/2 + 1/2)))$

using *d-gt-0* **by** (*auto simp: diamond-x-def diamond-y-def dest: fun-cong [where x = 1/2]*)

lemma *neqs2*:

shows $(\lambda x. (diamond\text{-}x\ x, diamond\text{-}y (diamond\text{-}x\ x))) \neq (\lambda x. ((2 * x - 1) * diamond\text{-}y (diamond\text{-}x\ 1), diamond\text{-}x\ 1))$

and $(\lambda x. (diamond\text{-}x\ x, - diamond\text{-}y (diamond\text{-}x\ x))) \neq (\lambda x. ((2 * x - 1) * diamond\text{-}y (diamond\text{-}x\ 0), diamond\text{-}x\ 0))$

using *d-gt-0* **by** (*auto simp: diamond-x-def diamond-y-def dest: fun-cong [where x = 1]*)

```

lemma diamond-cube-is-only-horizontal-div-of-rot:
  shows only-horizontal-division (boundary (diamond-cube)) {rot-diamond-cube}
  unfolding only-horizontal-division-def
proof (rule exI [of - {}], simp, intro conjI ballI)
  show finite (boundary diamond-cube)
    by (simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
  show boundary-chain (boundary diamond-cube)
    by (simp add: two-cube-boundary-is-boundary)
  show  $\bigwedge x. x \in \text{boundary diamond-cube} \implies \text{case } x \text{ of } (k, x) \implies \text{valid-path } x$ 
    using diamond-cube-boundary-valid by blast
  let  $\mathcal{V} = (\text{boundary (diamond-cube)})$ 
  have  $0$ : finite  $\mathcal{V}$ 
    by (auto simp add: boundary-def horizontal-boundary-def vertical-boundary-def)
  have  $1$ : boundary-chain  $\mathcal{V}$  using two-cube-boundary-is-boundary by auto
  let  $\text{?subdiv} = \{\text{diamond-top-left-edge}, \text{diamond-top-right-edge}, \text{diamond-bot-left-edge},$ 
diamond-bot-right-edge}\}
  let  $\text{?pi} = \{(1::\text{int}, \lambda x. ((2 * x - 1) * \text{diamond-y (diamond-x } 0), \text{diamond-x } 0)),$ 
 $(- 1, \lambda x. ((2 * x - 1) * \text{diamond-y (diamond-x } 1), \text{diamond-x } 1))\}$ 
  let  $\text{?pj} = \{(-1::\text{int}, \lambda y. (\text{diamond-x } 0, (2 * y - 1) * \text{diamond-y (diamond-x}$ 
 $0))),$ 
 $(1, \lambda y. (\text{diamond-x } 1, (2 * y - 1) * \text{diamond-y (diamond-x } 1)))\}$ 
  let  $\text{?f} = \text{diamond-cube-boundary-to-subdiv}$ 
  have  $2$ : common-sudiv-exists (two-chain-vertical-boundary {rot-diamond-cube}
 $\mathcal{V}$ 
  unfolding common-sudiv-exists-def
proof (intro exI conjI)
  show chain-subdiv-chain (boundary (diamond-cube) -  $\text{?pj}$ )  $\text{?subdiv}$ 
    unfolding chain-subdiv-chain-character
  proof (intro exI conjI)
  have  $1$ : (boundary (diamond-cube)) -  $\text{?pj} = \{\text{diamond-top-edges}, \text{diamond-bot-edges}\}$ 
    apply (auto simp add: diamond-cube-boundary-explicit diamond-x-def dia-
mond-top-edges-def diamond-bot-edges-def)
    apply (metis (no-types, opaque-lifting) abs-zero cancel-comm-monoid-add-class.diff-cancel
diamond-x-def diamond-y-def diff-0 minus-diff-eq mult commute mult-zero-right prod.inject
neqs2)
    by (metis (no-types, opaque-lifting) cancel-comm-monoid-add-class.diff-cancel
d-gt-0 divide-eq-0-iff linorder-not-le mult commute mult-zero-right order-refl prod.sel(1)
zero-neq-numeral)
  then show  $\bigcup (\text{?f ' (boundary (diamond-cube) - ?pj) = ?subdiv$ 
    by (auto simp add: diamond-top-edges-def diamond-bot-edges-def dia-
mond-cube-boundary-to-subdiv-def)
  have chain-subdiv-path (reversepath ( $\lambda x. (\text{diamond-x } x, \text{diamond-y (diamond-x}$ 
 $x))))$ 
 $\{( - 1::\text{int}, \lambda x. (\text{diamond-x}(x/2), \text{diamond-x}(x/2) + d/2)),$ 
 $( - 1::\text{int}, \lambda x. (\text{diamond-x}(x/2 + 1/2), d/2 - \text{diamond-x}(x/2 +$ 
 $1/2))\}$ 
proof -
  let  $\text{?F} = \lambda x. (\text{diamond-x}(x/2 + 1/2), d/2 - \text{diamond-x}(x/2 + 1/2))$ 

```

```

let ?G = λx. (diamond-x(x/2), diamond-x(x/2) + d/2)
have dist: distinct [(-1, ?F), (-1, ?G)]
  using d-gt-0 by (auto simp: diamond-x-def diamond-y-def dest: fun-cong)
  have rj: rec-join [(-1, ?F), (-1, ?G)] = reversepath (λx. (diamond-x x,
diamond-y (diamond-x x)))
  using d-gt-0 by (auto simp add: subpath-def diamond-x-def diamond-y-def
joinpaths-def reversepath-def algebra-simps divide-simps)
  have pathstart ?F = pathfinish ?G
  using d-gt-0 by(auto simp add: subpath-def diamond-x-def diamond-y-def
pathfinish-def pathstart-def)
  then have val: valid-chain-list [(-1, ?F), (-1, ?G)]
    by (auto simp add: join-subpaths-middle)
  show ?thesis
    using chain-subdiv-path.I [OF dist rj val]
    by (simp add: insert-commute)
qed
moreover have chain-subdiv-path (λx. (diamond-x x, - diamond-y (diamond-x
x)))
  { (1::int, λx. (diamond-x(x/2), - diamond-x(x/2) - d/2)),
    (1::int, λx. (diamond-x(x/2 + 1/2), diamond-x(x/2 +
1/2) - d/2)) }
proof -
  let ?F = λx. (diamond-x(x/2), - diamond-x(x/2) - d/2)
  let ?G = λx. (diamond-x(x/2 + 1/2), diamond-x(x/2 + 1/2) - d/2)
  have dist: distinct [(1, ?F), (1, ?G)]
    using d-gt-0 by (auto simp: diamond-x-def diamond-y-def dest: fun-cong)
  have rj: rec-join [(1, ?F), (1, ?G)] = (λx. (diamond-x x, - diamond-y
(diamond-x x)))
    using d-gt-0 by (auto simp add: subpath-def diamond-x-def diamond-y-def
joinpaths-def algebra-simps divide-simps)
  have pathfinish ?F = pathstart ?G
    using d-gt-0 by(auto simp add: subpath-def diamond-x-def diamond-y-def
pathfinish-def pathstart-def)
  then have val: valid-chain-list [(1, ?F), (1, ?G)]
    by (auto simp add: join-subpaths-middle)
  show ?thesis
    using chain-subdiv-path.I [OF dist rj val] by simp
qed
ultimately show **:
  ∀ (k::int, γ) ∈ boundary (diamond-cube) - ?pj.
    if k = (1::int) then chain-subdiv-path γ (?f (k::int, γ))
    else chain-subdiv-path (reversepath γ) (?f (k::int, γ))
  ∀ p ∈ boundary (diamond-cube) - ?pj. ∀ p' ∈ boundary (diamond-cube) - ?pj.
    p ≠ p' → ?f p ∩ ?f p' = {}
  ∀ x ∈ boundary (diamond-cube) - ?pj. finite (?f x)
    by(simp-all add: diamond-cube-boundary-to-subdiv-def UNION-eq 1 dia-
mond-top-edges-def diamond-bot-edges-def diamond-bot-left-edge-def diamond-bot-right-edge-def
diamond-top-left-edge-def diamond-top-right-edge-def neqs1)
qed

```

```

have *:  $\bigwedge f. \bigcup (f \text{ ' } \{\text{rot-diamond-cube}\}) = f (\text{rot-diamond-cube})$  by auto
show chain-subdiv-chain (two-chain-vertical-boundary {rot-diamond-cube} -
?pi) ?subdiv
unfolding chain-subdiv-chain-character two-chain-vertical-boundary-def *
proof (intro exI conjI)
let ?f = rot-diamond-cube-boundary-to-subdiv
have 1: (vertical-boundary (rot-diamond-cube) - ?pi) = {diamond-left-edges,
diamond-right-edges}
apply (auto simp add: rot-diamond-cube-vertical-boundary-explicit dia-
mond-left-edges-def diamond-right-edges-def)
apply (metis (no-types, opaque-lifting) add.inverse-inverse add-diff-cancel-right'
diff-numeral-special(11) mult.left-neutral mult.right-neutral prod.inject neqs1(2) umi-
nus-add-conv-diff)
by (metis (no-types, opaque-lifting) diff-0 mult.left-neutral mult-minus-left
mult-zero-right prod.inject neqs1(2))
show  $\bigcup ( ?f \text{ ' } (\text{vertical-boundary (rot-diamond-cube) - ?pi})) = ?subdiv$ 
apply (simp add: rot-diamond-cube-boundary-to-subdiv-def 1 UNION-eq
subpath-def)
apply (auto simp add: set-eq-iff diamond-right-edges-def diamond-left-edges-def)
done
have chain-subdiv-path (reversepath ( $\lambda y. (- \text{diamond-y (diamond-x y), dia-$ 
mond-x y)))

$$\{(1, \lambda x. (\text{diamond-x}(x/2), - \text{diamond-x}(x/2) - d/2)),$$


$$(-1, \lambda x. (\text{diamond-x}(x/2), \text{diamond-x}(x/2) + d/2))\}$$

proof -
let ?F =  $\lambda x. (\text{diamond-x}(x/2), - \text{diamond-x}(x/2) - d/2)$ 
let ?G =  $\lambda x. (\text{diamond-x}(x/2), \text{diamond-x}(x/2) + d/2)$ 
have dist: distinct [(-1, ?G), (1::int, ?F)]
using d-gt-0 by simp
have rj: rec-join [(-1, ?G), (1::int, ?F)] = reversepath ( $\lambda y. (- \text{diamond-y}$ 
(diamond-x y), diamond-x y))
using d-gt-0 by (auto simp add: subpath-def diamond-x-def diamond-y-def
joinpaths-def reversepath-def algebra-simps divide-simps)
have pathstart ?G = pathstart ?F
using d-gt-0 by (auto simp add: subpath-def diamond-x-def diamond-y-def
pathstart-def)
then have val: valid-chain-list [(-1, ?G), (1::int, ?F)]
by (auto simp add: join-subpaths-middle)
show ?thesis
using chain-subdiv-path.I [OF dist rj val] by (simp add: insert-commute)
qed
moreover have chain-subdiv-path ( $\lambda y. (\text{diamond-y (diamond-x y), diamond-x}$ 
y))

$$\{(1, \lambda x. (\text{diamond-x}(x/2 + 1/2), \text{diamond-x}(x/2 + 1/2) -$$


$$d/2)),$$


$$(-1, \lambda x. (\text{diamond-x}(x/2 + 1/2), d/2 - \text{diamond-x}(x/2$$


$$+ 1/2)))\}$$

proof -
let ?F =  $\lambda x. (\text{diamond-x}(x/2 + 1/2), d/2 - \text{diamond-x}(x/2 + 1/2))$ 

```

```

let ?G = λx. (diamond-x(x/2 + 1/2), diamond-x(x/2 + 1/2) - d/2)
have dist: distinct [(1::int, ?G), (-1, ?F)]
  by simp
  have rj: rec-join [(1::int, ?G), (-1, ?F)] = (λy. (diamond-y (diamond-x
y), diamond-x y))
  using d-gt-0 by (auto simp add: subpath-def diamond-x-def diamond-y-def
joinpaths-def reversepath-def algebra-simps divide-simps)
  have pathfinish ?F = pathfinish ?G
  using d-gt-0 by(auto simp add: subpath-def diamond-x-def diamond-y-def
pathfinish-def pathstart-def)
  then have val: valid-chain-list [(1::int, ?G), (-1, ?F)]
  by (auto simp add: join-subpaths-middle)
  show ?thesis
  using chain-subdiv-path.I [OF dist rj val] by simp
qed
ultimately show ∀(k, γ)∈vertical-boundary (rot-diamond-cube) - ?pi.
  if k = 1 then chain-subdiv-path γ (?f (k, γ))
  else chain-subdiv-path (reversepath γ) (?f (k, γ))
  ∀p∈vertical-boundary (rot-diamond-cube) - ?pi.
  ∀p'∈vertical-boundary (rot-diamond-cube) - ?pi.
  p ≠ p' → ?f p ∩ ?f p' = {}
  ∀x∈vertical-boundary (rot-diamond-cube) - ?pi. finite (?f x)
by(auto simp add: rot-diamond-cube-boundary-to-subdiv-def 1 diamond-left-edges-def
diamond-right-edges-def diamond-bot-left-edge-def diamond-bot-right-edge-def
diamond-top-left-edge-def diamond-top-right-edge-def neqs1)
qed
show (∀(k::int, γ)∈?subdiv. valid-path γ)
  by (simp add: diamond-edges-are-valid prod.case-eq-if set-eq-iff)
show boundary-chain ?subdiv
  by (auto simp add: boundary-chain-def diamond-top-left-edge-def diamond-top-right-edge-def
diamond-bot-left-edge-def diamond-bot-right-edge-def)
show (∀(k, γ)∈?pi. point-path γ)
  using d-gt-0 by (auto simp add: point-path-def diamond-x-def diamond-y-def)
show (∀(k, γ)∈?pj. point-path γ)
  using d-gt-0 by (auto simp add: point-path-def diamond-x-def diamond-y-def)
qed
show common-sudiv-exists (two-chain-vertical-boundary {rot-diamond-cube}) (boundary
diamond-cube) ∨
  common-reparam-exists (boundary diamond-cube) (two-chain-vertical-boundary
{rot-diamond-cube})
  using 0 1 2 diamond-cube-boundary-valid by auto
qed

abbreviation rot-y t1 t2 ≡ (t1 - 1/2) / (2* diamond-y-gen (x-coord (rot-x t1
t2))) + 1/2

lemma rot-y-ivl:

```

```

assumes  $(0::real) \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1$ 
shows  $0 \leq \text{rot-}y \ x \ y \wedge \text{rot-}y \ x \ y \leq 1$ 
using assms
apply(auto simp add: x-coord-def diamond-y-gen-def algebra-simps divide-simps
linorder-class.not-le abs-if)
using mult-nonneg-nonneg apply fastforce
using dual-order.order-iff-strict apply fastforce
apply (sos ((((( $A < 0 * A < 1$ ) *  $R < 1$ ) + ((( $A <= 4 * (A < 1 * R < 1)$ ) * ( $R < 1/2$ 
*  $[1]^{\wedge}2$ )) + ((( $A <= 3 * (A < 0 * R < 1)$ ) * ( $R < 1/2 * [1]^{\wedge}2$ )) + (( $A <= 0 * (A <= 2$ 
* ( $A <= 3 * R < 1$ ))) * ( $R < 1 * [1]^{\wedge}2$ )))))) & (((( $A < 0 * A < 1$ ) *  $R < 1$ ) + ((( $A <= 4$ 
* ( $A < 1 * R < 1$ ) * ( $R < 1/3 * [1]^{\wedge}2$ )) + ((( $A <= 4 * (A < 0 * R < 1)$ ) * ( $R < 1/3 *$ 
 $[1]^{\wedge}2$ )) + (( $A <= 3 * (A <= 4 * R < 1)$ ) * ( $R < 1/3 * [1]^{\wedge}2$ ))))))))))
using assms(1) mult-left-le-one-le apply blast
using affine-ineq apply fastforce+
done

```

lemma *diamond-gen-eq-rot-diamond:*

```

assumes  $0 \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1$ 
shows (diamond-cube-gen  $(x, y)$ ) = (rot-diamond-cube (rot-y  $x \ y$ , rot-x  $x \ y$ ))
proof
show snd (diamond-cube-gen  $(x, y)$ ) = snd (rot-diamond-cube (rot-y  $x \ y$ , rot-x  $x \ y$ ))
apply(simp only: rot-diamond-cube-def diamond-eq-characterisations-fun)
by(auto simp add: diamond-cube-gen-def x-coord-def diamond-y-gen-def algebra-simps divide-simps)
show fst (diamond-cube-gen  $(x, y)$ ) = fst (rot-diamond-cube (rot-y  $x \ y$ , rot-x  $x \ y$ ))
using assms
apply(auto simp add: diamond-cube-gen-def rot-diamond-cube-def diamond-eq-characterisations-fun)
apply(auto simp add: x-coord-def diamond-y-gen-def algebra-simps divide-simps
abs-if split: if-split-asm)
apply sos+
done
qed

```

lemma *rot-diamond-eq-diamond-gen:*

```

assumes  $0 \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1$ 
shows (rot-diamond-cube  $(x, y)$ ) = (diamond-cube-gen (rot-x  $y \ x$ , rot-y  $y \ x$ ))
proof
show fst (rot-diamond-cube  $(x, y)$ ) = fst (diamond-cube-gen (rot-x  $y \ x$ , rot-y  $y \ x$ ))
apply(simp only: rot-diamond-cube-def diamond-eq-characterisations-fun)
by(auto simp add: diamond-cube-gen-def x-coord-def diamond-y-gen-def algebra-simps divide-simps)
show snd (rot-diamond-cube  $(x, y)$ ) = snd (diamond-cube-gen (rot-x  $y \ x$ , rot-y  $y \ x$ ))
using assms
apply(auto simp add: diamond-cube-gen-def rot-diamond-cube-def diamond-eq-characterisations-fun)
apply(auto simp add: x-coord-def diamond-y-gen-def algebra-simps divide-simps)

```


abs-if split: if-split-asm)

apply *sos+*

done

qed

lemma *rot-img-eq: cubeImage (diamond-cube-gen) = cubeImage (rot-diamond-cube)*

proof(*auto simp add: cubeImage-def image-def cbox-def real-pair-basis*)

show $\exists a \geq 0. a \leq 1 \wedge (\exists b \geq 0. b \leq 1 \wedge \text{diamond-cube-gen } (x, y) = \text{rot-diamond-cube } (a, b))$

if $0 \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1 \ (a, b) = \text{diamond-cube-gen } (x, y)$

for $a \ b \ x \ y$

proof (*intro exI conjI*)

let $?a = \text{rot-y } x \ y$

let $?b = \text{rot-x } x \ y$

show $0 \leq ?a \ ?a \leq 1$

using *rot-y-ivl that by blast+*

show $0 \leq ?b \ ?b \leq 1$

using *rot-x-ivl that by blast+*

show $\text{diamond-cube-gen } (x, y) = \text{rot-diamond-cube } (?a, ?b)$

using *that d-gt-0 diamond-gen-eq-rot-diamond by auto*

qed

show $\exists a \geq 0. a \leq 1 \wedge (\exists b \geq 0. b \leq 1 \wedge \text{rot-diamond-cube } (x, y) = \text{diamond-cube-gen } (a, b))$

if $0 \leq x \ x \leq 1 \ 0 \leq y \ y \leq 1 \ (a, b) = \text{rot-diamond-cube } (x, y)$ **for** $a \ b \ x \ y$

proof (*intro exI conjI*)

let $?a = \text{rot-x } y \ x$

let $?b = \text{rot-y } y \ x$

show $0 \leq ?a \ ?a \leq 1$

using *rot-x-ivl that by blast+*

show $0 \leq ?b \ ?b \leq 1$

using *rot-y-ivl that by blast+*

show $\text{rot-diamond-cube } (x, y) = \text{diamond-cube-gen } (?a, ?b)$

using *that d-gt-0 rot-diamond-eq-diamond-gen by auto*

qed

qed

lemma *rot-diamond-gen-div-diamond-gen:*

shows *gen-division (cubeImage (diamond-cube-gen)) (cubeImage ‘{rot-diamond-cube})*

using *rot-img-eq by(auto simp add: gen-division-def)*

lemma *rot-diamond-gen-div-diamond:*

shows *gen-division (cubeImage (diamond-cube)) (cubeImage ‘{rot-diamond-cube})*

using *rot-diamond-gen-div-diamond-gen*

by(*simp only: diamond-eq-characterisations-fun*)

lemma *GreenThm-diamond:*

assumes *analytically-valid (cubeImage (diamond-cube)) ($\lambda x. F \ x \cdot \ i$) j*

analytically-valid (cubeImage (diamond-cube)) ($\lambda x. F \ x \cdot \ j$) i

shows *integral (cubeImage (diamond-cube)) ($\lambda x. \text{partial-vector-derivative } (\lambda x. F$*

$x \cdot j) i x - \text{partial-vector-derivative } (\lambda x. F x \cdot i) j x =$
 $\text{one-chain-line-integral } F \{i, j\} (\text{boundary } (\text{diamond-cube}))$

proof –

- have** 0: $\forall (k, \gamma) \in \text{boundary } (\text{diamond-cube}). \text{valid-path } \gamma$
- using** *diamond-cube-boundary-valid d-gt-0* **by** *auto*
- have** $\wedge \text{twoCube}. \text{twoCube} \in \{\text{diamond-cube}\} \implies \text{typeI-twoCube } \text{twoCube}$
- using** *assms diamond-cube-is-type-I* **by** *auto*
- moreover have** *valid-two-chain {diamond-cube}*
- using** *assms(1) diamond-cube-valid-two-cube valid-two-chain-def* **by** *auto*
- moreover have** *gen-division (cubeImage (diamond-cube)) (cubeImage ‘ {diamond-cube})*
- by** *(simp add: gen-division-def)*
- moreover have** $(\forall \text{twoCube} \in \{\text{rot-diamond-cube}\}). \text{typeII-twoCube } \text{twoCube}$
- using** *assms rot-diamond-cube-is-type-II* **by** *auto*
- moreover have** *valid-two-chain {rot-diamond-cube}*
- using** *assms(1) rot-diamond-cube-valid-two-cube valid-two-chain-def* **by** *auto*
- moreover have** *gen-division (cubeImage (diamond-cube)) (cubeImage ‘ {rot-diamond-cube})*
- using** *rot-diamond-gen-div-diamond* **by** *auto*
- moreover have** 3: *only-vertical-division (boundary (diamond-cube)) {diamond-cube}*
- using** *twoChainVertDiv-of-itself[of {diamond-cube}] diamond-cube-boundary-valid*

assms

- by** *(auto simp add: two-chain-boundary-def)*
- moreover have** 4: $\forall \text{twoC} \in \{\text{diamond-cube}\}. \text{analytically-valid } (\text{cubeImage } \text{twoC})$
- $(\lambda x. F x \cdot i) j$
- using** *assms*
- by** *fastforce*
- moreover have** 5: $\forall \text{twoC} \in \{\text{rot-diamond-cube}\}. \text{analytically-valid } (\text{cubeImage } \text{twoC})$
- $(\lambda x. F x \cdot j) i$
- using** *assms diamond-eq-characterisations-fun rot-img-eq* **by** *auto*
- ultimately show** *?thesis*
- using** *green-typeI-typeII-chain.GreenThm-typeI-typeII-divisible-region-finite-holes[of (cubeImage (diamond-cube)) i j F {diamond-cube} {rot-diamond-cube}, OF - 0 3 diamond-cube-is-only-horizontal-div-of-rot -]*
- R2-axioms*
- by** *(auto simp add: green-typeI-typeII-chain-def green-typeI-chain-def green-typeII-chain-def green-typeI-chain-axioms-def green-typeII-chain-axioms-def)*

qed
end
end