

Gauss Sums and the Pólya–Vinogradov Inequality

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Abstract

This article provides a full formalisation of Chapter 8 of Apostol’s *Introduction to Analytic Number Theory* [1]. Subjects that are covered are:

- periodic arithmetic functions and their finite Fourier series
- (generalised) Ramanujan sums
- Gauss sums and separable characters
- induced moduli and primitive characters
- the Pólya–Vinogradov inequality

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1 Auxiliary material

theory *Gauss-Sums-Auxiliary*

imports

Dirichlet-L.Dirichlet-Characters

Dirichlet-Series.Moebius-Mu

Dirichlet-Series.More-Totient

begin

1.1 Various facts

lemma *sum-div-reduce*:

fixes $d :: \text{nat}$ **and** $f :: \text{nat} \Rightarrow \text{complex}$

assumes $d \text{ dvd } k \ d > 0$

shows $(\sum n \mid n \in \{1..k\} \wedge d \text{ dvd } n. f \ n) = (\sum c \in \{1..k \text{ div } d\}. f \ (c*d))$

<proof>

lemma *prod-div-sub*:

fixes $f :: \text{nat} \Rightarrow \text{complex}$

assumes $\text{finite } A \ B \subseteq A \ \forall b \in B. f \ b \neq 0$

shows $(\prod i \in A - B. f \ i) = ((\prod i \in A. f \ i) \text{ div } (\prod i \in B. f \ i))$

<proof>

lemma *linear-gcd*:

fixes $a \ b \ c \ d :: \text{nat}$

assumes $a > 0 \ b > 0 \ c > 0 \ d > 0$

assumes $\text{coprime } a \ c \ \text{coprime } b \ d$

shows $\text{gcd } (a*b) \ (c*d) = (\text{gcd } a \ d) * (\text{gcd } b \ c)$

<proof>

lemma *reindex-product-bij*:

fixes $a \ b \ m \ k :: \text{nat}$

defines $S \equiv \{(d1,d2). d1 \text{ dvd } \text{gcd } a \ m \wedge d2 \text{ dvd } \text{gcd } k \ b\}$

defines $T \equiv \{d. d \text{ dvd } (\text{gcd } a \ m) * (\text{gcd } k \ b)\}$

defines $f \equiv (\lambda(d1,d2). d1 * d2)$

assumes $\text{coprime } a \ k$

shows *bij-betw* $f \ S \ T$

<proof>

lemma *p-div-set*:

shows $\{p. p \in \text{prime-factors } a \wedge \neg p \text{ dvd } N\} =$

$(\{p. p \in \text{prime-factors } (a*N)\} - \{p. p \in \text{prime-factors } N\})$

(is $?A = ?B$)

<proof>

lemma *coprime-iff-prime-factors-disjoint*:

fixes $x \ y :: 'a :: \text{factorial-semiring}$

assumes $x \neq 0 \ y \neq 0$

shows $\text{coprime } x \ y \longleftrightarrow \text{prime-factors } x \cap \text{prime-factors } y = \{\}$

<proof>

lemma *coprime-cong-prime-factors*:
fixes $x\ y :: 'a :: \text{factorial-semiring-gcd}$
assumes $x \neq 0\ y \neq 0\ x' \neq 0\ y' \neq 0$
assumes $\text{prime-factors } x = \text{prime-factors } x'$
assumes $\text{prime-factors } y = \text{prime-factors } y'$
shows $\text{coprime } x\ y \longleftrightarrow \text{coprime } x'\ y'$
 $\langle \text{proof} \rangle$

lemma *moebius-prod-not-coprime*:
assumes $\neg \text{coprime } N\ d$
shows $\text{moebius-mu } (N*d) = 0$
 $\langle \text{proof} \rangle$

Theorem 2.18

lemma *sum-divisors-moebius-mu-times-multiplicative*:
fixes $f :: \text{nat} \Rightarrow 'a :: \{\text{comm-ring-1}\}$
assumes *multiplicative-function* f **and** $n > 0$
shows $(\sum d \mid d\ \text{dvd}\ n. \text{moebius-mu } d * f\ d) = (\prod_{p \in \text{prime-factors } n}. 1 - f\ p)$
 $\langle \text{proof} \rangle$

lemma *multiplicative-ind-coprime [intro]*: *multiplicative-function* $(\text{ind } (\text{coprime } N))$
 $\langle \text{proof} \rangle$

lemma *sum-divisors-moebius-mu-times-multiplicative-revisited*:
fixes $f :: \text{nat} \Rightarrow 'a :: \{\text{comm-ring-1}\}$
assumes *multiplicative-function* f $n > 0\ N > 0$
shows $(\sum d \mid d\ \text{dvd}\ n \wedge \text{coprime } N\ d. \text{moebius-mu } d * f\ d) =$
 $(\prod_{p \in \{p. p \in \text{prime-factors } n \wedge \neg (p\ \text{dvd}\ N)\}}. 1 - f\ p)$
 $\langle \text{proof} \rangle$

1.2 Neutral element of the Dirichlet product

definition *dirichlet-prod-neutral* $n = (\text{if } n = 1 \text{ then } 1 \text{ else } 0)$ **for** $n :: \text{nat}$

lemma *dirichlet-prod-neutral-intro*:
fixes $S :: \text{nat} \Rightarrow \text{complex}$ **and** $f :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{complex}$
defines $S \equiv (\lambda(n::\text{nat}). (\sum k \mid k \in \{1..n\} \wedge \text{coprime } k\ n. (f\ k\ n)))$
shows $S(n) = (\sum k \in \{1..n\}. f\ k\ n * \text{dirichlet-prod-neutral } (\text{gcd } k\ n))$
 $\langle \text{proof} \rangle$

lemma *dirichlet-prod-neutral-right-neutral*:
 $\text{dirichlet-prod } f\ \text{dirichlet-prod-neutral } n = f\ n$ **if** $n > 0$ **for** $f :: \text{nat} \Rightarrow \text{complex}$
and n
 $\langle \text{proof} \rangle$

lemma *dirichlet-prod-neutral-left-neutral*:
 $\text{dirichlet-prod } \text{dirichlet-prod-neutral } f\ n = f\ n$
if $n > 0$ **for** $f :: \text{nat} \Rightarrow \text{complex}$ **and** n

<proof>

corollary *I-right-neutral-0:*

fixes $f :: \text{nat} \Rightarrow \text{complex}$

assumes $f\ 0 = 0$

shows $\text{dirichlet-prod } f \text{ dirichlet-prod-neutral } n = f\ n$

<proof>

1.3 Multiplicative functions

lemma *mult-id: multiplicative-function id*

<proof>

lemma *mult-moebius: multiplicative-function moebius-mu*

<proof>

lemma *mult-of-nat: multiplicative-function of-nat*

<proof>

lemma *mult-of-nat-c: completely-multiplicative-function of-nat*

<proof>

lemma *completely-multiplicative-nonzero:*

fixes $f :: \text{nat} \Rightarrow \text{complex}$

assumes *completely-multiplicative-function* f

$d \neq 0$

$\bigwedge p. \text{prime } p \implies f(p) \neq 0$

shows $f(d) \neq 0$

<proof>

lemma *multipl-div:*

fixes $m\ k\ d1\ d2 :: \text{nat}$ **and** $f :: \text{nat} \Rightarrow \text{complex}$

assumes *multiplicative-function* f $d1\ \text{dvd}\ m\ d2\ \text{dvd}\ k$ *coprime* $m\ k$

shows $f\ ((m*k)\ \text{div}\ (d1*d2)) = f(m\ \text{div}\ d1) * f(k\ \text{div}\ d2)$

<proof>

lemma *multipl-div-mono:*

fixes $m\ k\ d :: \text{nat}$ **and** $f :: \text{nat} \Rightarrow \text{complex}$

assumes *completely-multiplicative-function* f

$d\ \text{dvd}\ k\ d > 0$

$\bigwedge p. \text{prime } p \implies f(p) \neq 0$

shows $f\ (k\ \text{div}\ d) = f(k)\ \text{div}\ f(d)$

<proof>

lemma *comp-to-mult: completely-multiplicative-function* $f \implies$
multiplicative-function f

<proof>

end

2 Periodic arithmetic functions

theory *Periodic-Arithmetic*

imports

Complex-Main

HOL-Number-Theory.Cong

begin

definition

periodic-arithmetic $f\ k = (\forall n. f\ (n+k) = f\ n)$
for $n :: \text{int}$ **and** $k :: \text{nat}$ **and** $f :: \text{nat} \Rightarrow \text{complex}$

lemma *const-periodic-arithmetic*: *periodic-arithmetic* $(\lambda x. y)\ k$
<proof>

lemma *add-periodic-arithmetic*:

fixes $f\ g :: \text{nat} \Rightarrow \text{complex}$
assumes *periodic-arithmetic* $f\ k$
assumes *periodic-arithmetic* $g\ k$
shows *periodic-arithmetic* $(\lambda n. f\ n + g\ n)\ k$
<proof>

lemma *mult-periodic-arithmetic*:

fixes $f\ g :: \text{nat} \Rightarrow \text{complex}$
assumes *periodic-arithmetic* $f\ k$
assumes *periodic-arithmetic* $g\ k$
shows *periodic-arithmetic* $(\lambda n. f\ n * g\ n)\ k$
<proof>

lemma *scalar-mult-periodic-arithmetic*:

fixes $f :: \text{nat} \Rightarrow \text{complex}$ **and** $a :: \text{complex}$
assumes *periodic-arithmetic* $f\ k$
shows *periodic-arithmetic* $(\lambda n. a * f\ n)\ k$
<proof>

lemma *fin-sum-periodic-arithmetic-set*:

fixes $f\ g :: \text{nat} \Rightarrow \text{complex}$
assumes $\forall i \in A. \text{periodic-arithmetic}\ (h\ i)\ k$
shows *periodic-arithmetic* $(\lambda n. \sum i \in A. h\ i\ n)\ k$
<proof>

lemma *mult-period*:

assumes *periodic-arithmetic* $g\ k$
shows *periodic-arithmetic* $g\ (k*q)$
<proof>

lemma *unique-periodic-arithmetic-extension*:

assumes $k > 0$
assumes $\forall j < k. g\ j = h\ j$

assumes *periodic-arithmetic g k and periodic-arithmetic h k*
shows $g\ i = h\ i$
 $\langle proof \rangle$

lemma *periodic-arithmetic-sum-periodic-arithmetic:*
assumes *periodic-arithmetic f k*
shows $(\sum l \in \{m..n\}. f\ l) = (\sum l \in \{m+k..n+k\}. f\ l)$
 $\langle proof \rangle$

lemma *mod-periodic-arithmetic:*
fixes $n\ m :: nat$
assumes *periodic-arithmetic f k*
assumes $n\ mod\ k = m\ mod\ k$
shows $f\ n = f\ m$
 $\langle proof \rangle$

lemma *cong-periodic-arithmetic:*
assumes *periodic-arithmetic f k [a = b] (mod k)*
shows $f\ a = f\ b$
 $\langle proof \rangle$

lemma *cong-nat-imp-eq:*
fixes $m :: nat$
assumes $m > 0\ x \in \{a..<a+m\}\ y \in \{a..<a+m\}\ [x = y]\ (mod\ m)$
shows $x = y$
 $\langle proof \rangle$

lemma *inj-on-mod-nat:*
fixes $m :: nat$
assumes $m > 0$
shows *inj-on* $(\lambda x. x\ mod\ m)\ \{a..<a+m\}$
 $\langle proof \rangle$

lemma *bij-betw-mod-nat-atLeastLessThan:*
fixes $k\ d :: nat$
assumes $k > 0$
defines $g \equiv (\lambda i. nat\ ((int\ i - int\ d)\ mod\ int\ k) + d)$
shows *bij-betw* $(\lambda i. i\ mod\ k)\ \{d..<d+k\}\ \{..<k\}$
 $\langle proof \rangle$

lemma *periodic-arithmetic-sum-periodic-arithmetic-shift:*
fixes $k\ d :: nat$
assumes *periodic-arithmetic f k k > 0 d > 0*
shows $(\sum l \in \{0..k-1\}. f\ l) = (\sum l \in \{d..d+k-1\}. f\ l)$
 $\langle proof \rangle$

lemma *self-bij-0-k:*
fixes $a\ k :: nat$
assumes *coprime a k [a*i = 1] (mod k) k > 0*

shows *bij-betw* $(\lambda r. r * a \bmod k) \{0..k-1\} \{0..k-1\}$
<proof>

lemma *periodic-arithmetic-homothecy*:

assumes *periodic-arithmetic* $f\ k$

shows *periodic-arithmetic* $(\lambda l. f\ (l * a))\ k$
<proof>

theorem *periodic-arithmetic-remove-homothecy*:

assumes *coprime* $a\ k$ *periodic-arithmetic* $f\ k\ k > 0$

shows $(\sum_{l=1..k} f\ l) = (\sum_{l=1..k} f\ (l * a))$
<proof>

end

theory *Complex-Roots-Of-Unity*

imports

HOL-Analysis.Analysis

Periodic-Arithmetic

begin

3 Complex roots of unity

definition

unity-root $k\ n = \text{cis}\ (2 * \pi * \text{of-int}\ n / \text{of-nat}\ k)$

lemma

unity-root-k-0 [*simp*]: *unity-root* $k\ 0 = 1$ **and**

unity-root-0-n [*simp*]: *unity-root* $0\ n = 1$

<proof>

lemma *unity-root-conv-exp*:

unity-root $k\ n = \text{exp}\ (\text{of-real}\ (2 * \pi * n / k) * i)$

<proof>

lemma *unity-root-mod*:

unity-root $k\ (n \bmod \text{int}\ k) = \text{unity-root}\ k\ n$

<proof>

lemma *unity-root-cong*:

assumes $[m = n] \bmod\ \text{int}\ k$

shows *unity-root* $k\ m = \text{unity-root}\ k\ n$

<proof>

lemma *unity-root-mod-nat*:

unity-root $k\ (\text{nat}\ (n \bmod \text{int}\ k)) = \text{unity-root}\ k\ n$

<proof>

lemma *unity-root-eqD*:

assumes *gr*: $k > 0$
assumes *eq*: $\text{unity-root } k \ i = \text{unity-root } k \ j$
shows $i \bmod k = j \bmod k$
 $\langle \text{proof} \rangle$

lemma *unity-root-eq-1-iff*:
fixes $k \ n :: \text{nat}$
assumes $k > 0$
shows $\text{unity-root } k \ n = 1 \iff k \ \text{dvd} \ n$
 $\langle \text{proof} \rangle$

lemma *unity-root-pow*: $\text{unity-root } k \ n^{\wedge} m = \text{unity-root } k \ (n * m)$
 $\langle \text{proof} \rangle$

lemma *unity-root-add*: $\text{unity-root } k \ (m + n) = \text{unity-root } k \ m * \text{unity-root } k \ n$
 $\langle \text{proof} \rangle$

lemma *unity-root-uminus*: $\text{unity-root } k \ (-m) = \text{cnj} \ (\text{unity-root } k \ m)$
 $\langle \text{proof} \rangle$

lemma *inverse-unity-root*: $\text{inverse} \ (\text{unity-root } k \ m) = \text{cnj} \ (\text{unity-root } k \ m)$
 $\langle \text{proof} \rangle$

lemma *unity-root-diff*: $\text{unity-root } k \ (m - n) = \text{unity-root } k \ m * \text{cnj} \ (\text{unity-root } k \ n)$
 $\langle \text{proof} \rangle$

lemma *unity-root-eq-1-iff-int*:
fixes $k :: \text{nat}$ **and** $n :: \text{int}$
assumes $k > 0$
shows $\text{unity-root } k \ n = 1 \iff k \ \text{dvd} \ n$
 $\langle \text{proof} \rangle$

lemma *unity-root-eq-1 [simp]*: $\text{int } k \ \text{dvd} \ n \implies \text{unity-root } k \ n = 1$
 $\langle \text{proof} \rangle$

lemma *unity-periodic-arithmetic*:
 $\text{periodic-arithmetic} \ (\text{unity-root } k) \ k$
 $\langle \text{proof} \rangle$

lemma *unity-periodic-arithmetic-mult*:
 $\text{periodic-arithmetic} \ (\lambda n. \text{unity-root } k \ (m * \text{int } n)) \ k$
 $\langle \text{proof} \rangle$

lemma *unity-root-periodic-arithmetic-mult-minus*:
shows $\text{periodic-arithmetic} \ (\lambda i. \text{unity-root } k \ (-\text{int } i * \text{int } m)) \ k$
 $\langle \text{proof} \rangle$

lemma *unity-div*:

fixes $a :: \text{int}$ **and** $d :: \text{nat}$
assumes $d \text{ dvd } k$
shows $\text{unity-root } k (a*d) = \text{unity-root } (k \text{ div } d) a$
 $\langle \text{proof} \rangle$

lemma *unity-div-num*:
assumes $k > 0$ $d > 0$ $d \text{ dvd } k$
shows $\text{unity-root } k (x * (k \text{ div } d)) = \text{unity-root } d x$
 $\langle \text{proof} \rangle$

4 Geometric sums of roots of unity

Apostol calls these ‘geometric sums’, which is a bit too generic. We therefore decided to refer to them as ‘sums of roots of unity’.

definition $\text{unity-root-sum } k n = (\sum_{m < k} \text{unity-root } k (n * \text{of-nat } m))$

lemma *unity-root-sum-0-left* [simp]: $\text{unity-root-sum } 0 n = 0$ **and**
unity-root-sum-0-right [simp]: $k > 0 \implies \text{unity-root-sum } k 0 = k$
 $\langle \text{proof} \rangle$

Theorem 8.1

theorem *unity-root-sum*:
fixes $k :: \text{nat}$ **and** $n :: \text{int}$
assumes $gr: k \geq 1$
shows $k \text{ dvd } n \implies \text{unity-root-sum } k n = k$
and $\neg k \text{ dvd } n \implies \text{unity-root-sum } k n = 0$
 $\langle \text{proof} \rangle$

corollary *unity-root-sum-periodic-arithmetic*:
periodic-arithmetic $(\text{unity-root-sum } k) k$
 $\langle \text{proof} \rangle$

lemma *unity-root-sum-nonzero-iff*:
fixes $r :: \text{int}$
assumes $k \geq 1$ **and** $r \in \{-k < .. < k\}$
shows $\text{unity-root-sum } k r \neq 0 \iff r = 0$
 $\langle \text{proof} \rangle$

end

5 Finite Fourier series

theory *Finite-Fourier-Series*
imports
Polynomial-Interpolation.Lagrange-Interpolation
Complex-Roots-Of-Unity
begin

5.1 Auxiliary facts

lemma *lagrange-exists*:

assumes *d*: *distinct* (*map fst zs-ws*)

defines *e*: (*p* :: *complex poly*) \equiv *lagrange-interpolation-poly* *zs-ws*

shows *degree* *p* \leq (*length* *zs-ws*) - 1

$(\forall x y. (x,y) \in \text{set } zs\text{-}ws \longrightarrow \text{poly } p \ x = y)$

<proof>

lemma *lagrange-unique*:

assumes *o*: *length* *zs-ws* > 0

assumes *d*: *distinct* (*map fst zs-ws*)

assumes 1: *degree* (*p1* :: *complex poly*) \leq (*length* *zs-ws*) - 1 \wedge

$(\forall x y. (x,y) \in \text{set } zs\text{-}ws \longrightarrow \text{poly } p1 \ x = y)$

assumes 2: *degree* (*p2* :: *complex poly*) \leq (*length* *zs-ws*) - 1 \wedge

$(\forall x y. (x,y) \in \text{set } zs\text{-}ws \longrightarrow \text{poly } p2 \ x = y)$

shows *p1* = *p2*

<proof>

Theorem 8.2

corollary *lagrange*:

assumes *length* *zs-ws* > 0 *distinct* (*map fst zs-ws*)

shows $(\exists! (p :: \text{complex poly}).$

degree *p* \leq *length* *zs-ws* - 1 \wedge

$(\forall x y. (x, y) \in \text{set } zs\text{-}ws \longrightarrow \text{poly } p \ x = y))$

<proof>

lemma *poly-altdef'*:

assumes *gr*: *k* \geq *degree* *p*

shows *poly* *p* (*z*::*complex*) = $(\sum_{i \leq k}. \text{coeff } p \ i * z^i)$

<proof>

5.2 Definition and uniqueness

definition *finite-fourier-poly* :: *complex list* \Rightarrow *complex poly* **where**

finite-fourier-poly *ws* =

(*let* *k* = *length* *ws*

in *poly-of-list* [*1 / k* * $(\sum_{m < k}. \text{ws} \ ! \ m * \text{unity-root } k \ (-n * m))$]. *n* \leftarrow [*0..<k*]])

lemma *degree-poly-of-list-le*: *degree* (*poly-of-list* *ws*) \leq *length* *ws* - 1

<proof>

lemma *degree-finite-fourier-poly*: *degree* (*finite-fourier-poly* *ws*) \leq *length* *ws* - 1

<proof>

lemma *coeff-finite-fourier-poly*:

assumes *n* < *length* *ws*

defines *k* \equiv *length* *ws*

shows *coeff* (*finite-fourier-poly* *ws*) *n* =

$(1/k) * (\sum_{m < k}. \text{ws} \ ! \ m * \text{unity-root } k \ (-n * m))$

<proof>

lemma *poly-finite-fourier-poly*:

fixes $m :: \text{int}$ **and** ws

defines $k \equiv \text{length } ws$

assumes $m \in \{0..<k\}$

assumes $m < \text{length } ws$

shows $\text{poly } (\text{finite-fourier-poly } ws) (\text{unity-root } k \ m) = ws \ ! \ (\text{nat } m)$

<proof>

Theorem 8.3

theorem *finite-fourier-poly-unique*:

assumes $\text{length } ws > 0$

defines $k \equiv \text{length } ws$

assumes $(\text{degree } p \leq k - 1)$

assumes $(\forall m \leq k-1. (ws \ ! \ m) = \text{poly } p (\text{unity-root } k \ m))$

shows $p = \text{finite-fourier-poly } ws$

<proof>

The following alternative formulation returns a coefficient

definition *finite-fourier-poly'* :: $(\text{nat} \Rightarrow \text{complex}) \Rightarrow \text{nat} \Rightarrow \text{complex poly}$ **where**

finite-fourier-poly' $ws \ k =$

$(\text{poly-of-list } [1 / k * (\sum m < k. (ws \ m) * \text{unity-root } k \ (-n * m)). n \leftarrow [0..<k]])$

lemma *finite-fourier-poly'-conv-finite-fourier-poly*:

finite-fourier-poly' $ws \ k = \text{finite-fourier-poly } [ws \ n. n \leftarrow [0..<k]]$

<proof>

lemma *coeff-finite-fourier-poly'*:

assumes $n < k$

shows $\text{coeff } (\text{finite-fourier-poly}' \ ws \ k) \ n =$

$(1/k) * (\sum m < k. (ws \ m) * \text{unity-root } k \ (-n * m))$

<proof>

lemma *degree-finite-fourier-poly'*: $\text{degree } (\text{finite-fourier-poly}' \ ws \ k) \leq k - 1$

<proof>

lemma *poly-finite-fourier-poly'*:

fixes $m :: \text{int}$ **and** k

assumes $m \in \{0..<k\}$

shows $\text{poly } (\text{finite-fourier-poly}' \ ws \ k) (\text{unity-root } k \ m) = ws \ (\text{nat } m)$

<proof>

lemma *finite-fourier-poly'-unique*:

assumes $k > 0$

assumes $\text{degree } p \leq k - 1$

assumes $\forall m \leq k-1. ws \ m = \text{poly } p (\text{unity-root } k \ m)$

shows $p = \text{finite-fourier-poly}' \ ws \ k$

<proof>

lemma *fourier-unity-root*:
fixes $k :: \text{nat}$
assumes $k > 0$
shows $\text{poly } (\text{finite-fourier-poly}' f k) (\text{unity-root } k m) =$
 $(\sum n < k. 1/k * (\sum m < k. (f m) * \text{unity-root } k (-n * m)) * \text{unity-root } k (m * n))$
 $\langle \text{proof} \rangle$

5.3 Expansion of an arithmetical function

Theorem 8.4

theorem *fourier-expansion-periodic-arithmetic*:
assumes $k > 0$
assumes *periodic-arithmetic* $f k$
defines $g \equiv (\lambda n. (1 / k) * (\sum m < k. f m * \text{unity-root } k (-n * m)))$
shows *periodic-arithmetic* $g k$
and $f m = (\sum n < k. g n * \text{unity-root } k (m * n))$
 $\langle \text{proof} \rangle$

theorem *fourier-expansion-periodic-arithmetic-unique*:
fixes $f g :: \text{nat} \Rightarrow \text{complex}$
assumes $k > 0$
assumes *periodic-arithmetic* $f k$ **and** *periodic-arithmetic* $g k$
assumes $\bigwedge m. m < k \implies f m = (\sum n < k. g n * \text{unity-root } k (\text{int } (m * n)))$
shows $g n = (1 / k) * (\sum m < k. f m * \text{unity-root } k (-n * m))$
 $\langle \text{proof} \rangle$

end

6 Ramanujan sums

theory *Ramanujan-Sums*
imports
Dirichlet-Series.Moebius-Mu
Gauss-Sums-Auxiliary
Finite-Fourier-Series
begin

6.1 Basic sums

definition *ramanujan-sum* $:: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{complex}$
where $\text{ramanujan-sum } k n = (\sum m \mid m \in \{1..k\} \wedge \text{coprime } m k. \text{unity-root } k (m * n))$

notation *ramanujan-sum* (c)

lemma *ramanujan-sum-0-n* [*simp*]: $c \ 0 \ n = 0$
 $\langle \text{proof} \rangle$

lemma *sum-coprime-conv-dirichlet-prod-moebius-mu*:
fixes $F S :: \text{nat} \Rightarrow \text{complex}$ **and** $f :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{complex}$
defines $F \equiv (\lambda n. (\sum k \in \{1..n\}. f k n))$
defines $S \equiv (\lambda n. (\sum k \mid k \in \{1..n\} \wedge \text{coprime } k n . f k n))$
assumes $\bigwedge a b d. d \text{ dvd } a \implies d \text{ dvd } b \implies f (a \text{ div } d) (b \text{ div } d) = f a b$
shows $S n = \text{dirichlet-prod moebius-mu } F n$
 $\langle \text{proof} \rangle$

lemma *dirichlet-prod-neutral-sum*:
 $\text{dirichlet-prod-neutral } n = (\sum k = 1..n. \text{unity-root } n k)$ **for** $n :: \text{nat}$
 $\langle \text{proof} \rangle$

lemma *moebius-coprime-sum*:
 $\text{moebius-mu } n = (\sum k \mid k \in \{1..n\} \wedge \text{coprime } k n . \text{unity-root } n (\text{int } k))$
 $\langle \text{proof} \rangle$

corollary *ramanujan-sum-1-right* [*simp*]: $c k (\text{Suc } 0) = \text{moebius-mu } k$
 $\langle \text{proof} \rangle$

lemma *ramanujan-sum-dvd-eq-totient*:
assumes $k \text{ dvd } n$
shows $c k n = \text{totient } k$
 $\langle \text{proof} \rangle$

6.2 Generalised sums

definition *gen-ramanujan-sum* :: $(\text{nat} \Rightarrow \text{complex}) \Rightarrow (\text{nat} \Rightarrow \text{complex}) \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{complex}$ **where**
 $\text{gen-ramanujan-sum } f g = (\lambda k n. \sum d \mid d \text{ dvd } \text{gcd } n k. f d * g (k \text{ div } d))$

notation *gen-ramanujan-sum* (s)

lemma *gen-ramanujan-sum-k-1*: $s f g k 1 = f 1 * g k$
 $\langle \text{proof} \rangle$

lemma *gen-ramanujan-sum-1-n*: $s f g 1 n = f 1 * g 1$
 $\langle \text{proof} \rangle$

lemma *gen-ramanujan-sum-periodic*: *periodic-arithmetic* ($s f g k$) k
 $\langle \text{proof} \rangle$

Theorem 8.5

theorem *gen-ramanujan-sum-fourier-expansion*:
fixes $f g :: \text{nat} \Rightarrow \text{complex}$ **and** $a :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{complex}$
assumes $k > 0$
defines $a \equiv (\lambda k m. (1/k) * (\sum d \mid d \text{ dvd } (\text{gcd } m k). g d * f (k \text{ div } d) * d))$
shows $s f g k n = (\sum m \leq k-1. a k m * \text{unity-root } k (m*n))$
 $\langle \text{proof} \rangle$

Theorem 8.6

theorem *ramanujan-sum-dirichlet-form*:

fixes $k\ n :: \text{nat}$

assumes $k > 0$

shows $c\ k\ n = (\sum d \mid d\ \text{dvd}\ \text{gcd}\ n\ k.\ d * \text{moebius-mu}\ (k\ \text{div}\ d))$

<proof>

corollary *ramanujan-sum-conv-gen-ramanujan-sum*:

$k > 0 \implies c\ k\ n = s\ \text{id}\ \text{moebius-mu}\ k\ n$

<proof>

Theorem 8.7

theorem *gen-ramanujan-sum-distrib*:

fixes $f\ g :: \text{nat} \Rightarrow \text{complex}$

assumes $a > 0\ b > 0\ m > 0\ k > 0$

assumes *coprime* $a\ k$ *coprime* $b\ m$ *coprime* $k\ m$

assumes *multiplicative-function* f **and**

multiplicative-function g

shows $s\ f\ g\ (m*k)\ (a*b) = s\ f\ g\ m\ a * s\ f\ g\ k\ b$

<proof>

corollary *gen-ramanujan-sum-distrib-right*:

fixes $f\ g :: \text{nat} \Rightarrow \text{complex}$

assumes $a > 0$ **and** $b > 0$ **and** $m > 0$

assumes *coprime* $b\ m$

assumes *multiplicative-function* f **and**

multiplicative-function g

shows $s\ f\ g\ m\ (a * b) = s\ f\ g\ m\ a$

<proof>

corollary *gen-ramanujan-sum-distrib-left*:

fixes $f\ g :: \text{nat} \Rightarrow \text{complex}$

assumes $a > 0$ **and** $k > 0$ **and** $m > 0$

assumes *coprime* $a\ k$ **and** *coprime* $k\ m$

assumes *multiplicative-function* f **and**

multiplicative-function g

shows $s\ f\ g\ (m*k)\ a = s\ f\ g\ m\ a * g\ k$

<proof>

corollary *ramanujan-sum-distrib*:

assumes $a > 0$ **and** $k > 0$ **and** $m > 0$ **and** $b > 0$

assumes *coprime* $a\ k$ *coprime* $b\ m$ *coprime* $m\ k$

shows $c\ (m*k)\ (a*b) = c\ m\ a * c\ k\ b$

<proof>

corollary *ramanujan-sum-distrib-right*:

assumes $a > 0$ **and** $k > 0$ **and** $m > 0$ **and** $b > 0$

assumes *coprime* $b\ m$

shows $c\ m\ (a*b) = c\ m\ a$

<proof>

corollary *ramanujan-sum-distrib-left:*

assumes $a > 0$ $k > 0$ $m > 0$

assumes *coprime* a k *coprime* m k

shows $c (m*k) a = c m a * \text{moebius-mu } k$

<proof>

lemma *dirichlet-prod-completely-multiplicative-left:*

fixes $f h :: \text{nat} \Rightarrow \text{complex}$ **and** $k :: \text{nat}$

defines $g \equiv (\lambda k. \text{moebius-mu } k * h k)$

defines $F \equiv \text{dirichlet-prod } f g$

assumes $k > 0$

assumes *completely-multiplicative-function* f
multiplicative-function h

assumes $\bigwedge p. \text{prime } p \implies f(p) \neq 0 \wedge f(p) \neq h(p)$

shows $F k = f k * (\prod_{p \in \text{prime-factors } k} 1 - h p / f p)$

<proof>

Theorem 8.8

theorem *gen-ramanujan-sum-dirichlet-expr:*

fixes $f h :: \text{nat} \Rightarrow \text{complex}$ **and** $n k :: \text{nat}$

defines $g \equiv (\lambda k. \text{moebius-mu } k * h k)$

defines $F \equiv \text{dirichlet-prod } f g$

defines $N \equiv k \text{ div gcd } n k$

assumes *completely-multiplicative-function* f
multiplicative-function h

assumes $\bigwedge p. \text{prime } p \implies f(p) \neq 0 \wedge f(p) \neq h(p)$

assumes $k > 0$ $n > 0$

shows $s f g k n = (F(k)*g(N)) \text{ div } (F(N))$

<proof>

lemma *totient-conv-moebius-mu-of-nat:*

of-nat (totient n) = dirichlet-prod moebius-mu of-nat n

<proof>

corollary *ramanujan-sum-k-n-dirichlet-expr:*

fixes $k n :: \text{nat}$

assumes $k > 0$ $n > 0$

shows $c k n = \text{of-nat (totient } k) * \text{moebius-mu (} k \text{ div gcd } n k) \text{ div of-nat (totient (} k \text{ div gcd } n k))$

<proof>

no-notation *ramanujan-sum (c)*

no-notation *gen-ramanujan-sum (s)*

end


```

theory Gauss-Sums
imports
  HOL-Algebra.Coset
  HOL-Real-Asymp.Real-Asymp
  Ramanujan-Sums
begin

```

7 Gauss sums

```

bundle vec-lambda-notation
begin
notation vec-lambda (binder  $\chi$  10)
end

```

```

bundle no-vec-lambda-notation
begin
no-notation vec-lambda (binder  $\chi$  10)
end

```

```

unbundle no-vec-lambda-notation

```

7.1 Definition and basic properties

```

context dcharacter
begin

```

```

lemma dir-periodic-arithmetic: periodic-arithmetic  $\chi$   $n$ 
  <proof>

```

```

definition gauss-sum  $k = (\sum m = 1..n . \chi(m) * \text{unity-root } n (m*k))$ 

```

```

lemma gauss-sum-periodic:
  periodic-arithmetic ( $\lambda n. \text{gauss-sum } n$ )  $n$ 
  <proof>

```

```

lemma ramanujan-sum-conv-gauss-sum:
  assumes  $\chi = \text{principal-dchar } n$ 
  shows ramanujan-sum  $n$   $k = \text{gauss-sum } k$ 
  <proof>

```

```

lemma cnj-mult-self:
  assumes coprime  $k$   $n$ 
  shows cnj ( $\chi$   $k$ ) *  $\chi$   $k = 1$ 
  <proof>

```

Theorem 8.9

```

theorem gauss-sum-reduction:

```

assumes *coprime k n*
shows $\text{gauss-sum } k = \text{cnj } (\chi \ k) * \text{gauss-sum } 1$
 ⟨*proof*⟩

The following variant takes an integer argument instead.

definition $\text{gauss-sum-int } k = (\sum_{m=1..n} \chi \ m * \text{unity-root } n \ (int \ m*k))$

sublocale *gauss-sum-int: periodic-fun-simple gauss-sum-int int n*
 ⟨*proof*⟩

lemma *gauss-sum-int-cong:*
assumes $[a = b] \ (mod \ int \ n)$
shows $\text{gauss-sum-int } a = \text{gauss-sum-int } b$
 ⟨*proof*⟩

lemma *gauss-sum-conv-gauss-sum-int:*
 $\text{gauss-sum } k = \text{gauss-sum-int } (int \ k)$
 ⟨*proof*⟩

lemma *gauss-sum-int-conv-gauss-sum:*
 $\text{gauss-sum-int } k = \text{gauss-sum } (nat \ (k \ mod \ n))$
 ⟨*proof*⟩

lemma *gauss-int-periodic: periodic-arithmetic gauss-sum-int n*
 ⟨*proof*⟩

proposition *dcharacter-fourier-expansion:*
 $\chi \ m = (\sum_{k=1..n} 1 / n * \text{gauss-sum-int } (-k) * \text{unity-root } n \ (m*k))$
 ⟨*proof*⟩

7.2 Separability

definition $\text{separable } k \longleftrightarrow \text{gauss-sum } k = \text{cnj } (\chi \ k) * \text{gauss-sum } 1$

corollary *gauss-coprime-separable:*
assumes *coprime k n*
shows *separable k*
 ⟨*proof*⟩

Theorem 8.10

theorem *global-separability-condition:*
 $(\forall n > 0. \text{separable } n) \longleftrightarrow (\forall k > 0. \neg \text{coprime } k \ n \longrightarrow \text{gauss-sum } k = 0)$
 ⟨*proof*⟩

lemma *of-real-moebius-mu [simp]: of-real (moebius-mu k) = moebius-mu k*
 ⟨*proof*⟩

corollary *principal-not-totally-separable:*
assumes $\chi = \text{principal-dchar } n$

shows $\neg(\forall k > 0. \text{separable } k)$
 ⟨proof⟩

Theorem 8.11

theorem *gauss-sum-1-mod-square-eq-k*:
assumes $(\forall k. k > 0 \longrightarrow \text{separable } k)$
shows $\text{norm } (\text{gauss-sum } 1) ^ 2 = \text{real } n$
 ⟨proof⟩

Theorem 8.12

theorem *gauss-sum-nonzero-noncoprime-necessary-condition*:
assumes $\text{gauss-sum } k \neq 0 \neg \text{coprime } k \ n \ k > 0$
defines $d \equiv n \ \text{div} \ \text{gcd } k \ n$
assumes $\text{coprime } a \ n \ [a = 1] \ (\text{mod } d)$
shows $d \ \text{dvd} \ n \ d < n \ \chi \ a = 1$
 ⟨proof⟩

7.3 Induced moduli and primitive characters

definition *induced-modulus* $d \longleftrightarrow d \ \text{dvd} \ n \wedge (\forall a. \text{coprime } a \ n \wedge [a = 1] \ (\text{mod } d) \longrightarrow \chi \ a = 1)$

lemma *induced-modulus-dvd*: *induced-modulus* $d \implies d \ \text{dvd} \ n$
 ⟨proof⟩

lemma *induced-modulusI* [*intro?*]:
 $d \ \text{dvd} \ n \implies (\bigwedge a. \text{coprime } a \ n \implies [a = 1] \ (\text{mod } d) \implies \chi \ a = 1) \implies \text{induced-modulus } d$
 ⟨proof⟩

lemma *induced-modulusD*: *induced-modulus* $d \implies \text{coprime } a \ n \implies [a = 1] \ (\text{mod } d) \implies \chi \ a = 1$
 ⟨proof⟩

lemma *zero-not-ind-mod*: $\neg \text{induced-modulus } 0$
 ⟨proof⟩

lemma *div-gcd-dvd1*: $(a :: 'a :: \text{semiring-gcd}) \ \text{div} \ \text{gcd } a \ b \ \text{dvd} \ a$
 ⟨proof⟩

lemma *div-gcd-dvd2*: $(b :: 'a :: \text{semiring-gcd}) \ \text{div} \ \text{gcd } a \ b \ \text{dvd} \ b$
 ⟨proof⟩

lemma *g-non-zero-ind-mod*:
assumes $\text{gauss-sum } k \neq 0 \neg \text{coprime } k \ n \ k > 0$
shows *induced-modulus* $(n \ \text{div} \ \text{gcd } k \ n)$
 ⟨proof⟩

lemma *induced-modulus-modulus*: *induced-modulus* n

<proof>

Theorem 8.13

theorem *one-induced-iff-principal:*
induced-modulus 1 \longleftrightarrow *$\chi =$ principal-dchar n*
<proof>

end

locale *primitive-dchar = dcharacter +*
assumes *no-induced-modulus:* $\neg(\exists d < n. \textit{induced-modulus } d)$

locale *nonprimitive-dchar = dcharacter +*
assumes *induced-modulus:* $\exists d < n. \textit{induced-modulus } d$

lemma (**in** *nonprimitive-dchar*) *nonprimitive:* $\neg \textit{primitive-dchar } n \ \chi$
<proof>

lemma (**in** *dcharacter*) *primitive-dchar-iff:*
primitive-dchar $n \ \chi \longleftrightarrow \neg(\exists d < n. \textit{induced-modulus } d)$
<proof>

lemma (**in** *residues-nat*) *principal-not-primitive:*
 $\neg \textit{primitive-dchar } n \ (\textit{principal-dchar } n)$
<proof>

lemma (**in** *dcharacter*) *not-primitive-imp-nonprimitive:*
assumes $\neg \textit{primitive-dchar } n \ \chi$
shows *nonprimitive-dchar* $n \ \chi$
<proof>

Theorem 8.14

theorem (**in** *dcharacter*) *prime-nonprincipal-is-primitive:*
assumes *prime* n
assumes $\chi \neq \textit{principal-dchar } n$
shows *primitive-dchar* $n \ \chi$
<proof>

Theorem 8.15

corollary (**in** *primitive-dchar*) *primitive-encoding:*
 $\forall k > 0. \neg \textit{coprime } k \ n \longrightarrow \textit{gauss-sum } k = 0$
 $\forall k > 0. \textit{separable } k$
 $\textit{norm } (\textit{gauss-sum } 1) \wedge 2 = n$
<proof>

Theorem 8.16

lemma (**in** *dcharacter*) *induced-modulus-altdef1:*

induced-modulus $d \longleftrightarrow$
 $d \text{ dvd } n \wedge (\forall a \ b. \text{ coprime } a \ n \wedge \text{ coprime } b \ n \wedge [a = b] \pmod{d} \longrightarrow \chi \ a = \chi \ b)$
 ⟨*proof*⟩

Exercise 8.4

lemma *induced-modulus-altdef2-lemma*:

fixes $n \ a \ d \ q :: \text{nat}$
defines $q \equiv (\prod p \mid \text{prime } p \wedge p \text{ dvd } n \wedge \neg (p \text{ dvd } a). \ p)$
defines $m \equiv a + q * d$
assumes $n > 0 \ \text{coprime } a \ d$
shows $[m = a] \pmod{d}$ **and** $\text{coprime } m \ n$
 ⟨*proof*⟩

Theorem 8.17

The case $d = 1$ is exactly the case described in *dcharacter* $?n \ ?\chi \implies \text{dcharacter.induced-modulus } ?n \ ?\chi \ 1 = (?\chi = \text{principal-dchar } ?n)$.

theorem (in *dcharacter*) *induced-modulus-altdef2*:

assumes $d \text{ dvd } n \ d \neq 1$
defines $\chi_1 \equiv \text{principal-dchar } n$
shows *induced-modulus* $d \longleftrightarrow (\exists \Phi. \text{dcharacter } d \ \Phi \wedge (\forall k. \chi \ k = \Phi \ k * \chi_1 \ k))$
 ⟨*proof*⟩

7.4 The conductor of a character

context *dcharacter*
begin

definition *conductor* = $\text{Min } \{d. \text{induced-modulus } d\}$

lemma *conductor-fin*: $\text{finite } \{d. \text{induced-modulus } d\}$
 ⟨*proof*⟩

lemma *conductor-induced*: $\text{induced-modulus } \text{conductor}$
 ⟨*proof*⟩

lemma *conductor-le-iff*: $\text{conductor} \leq a \longleftrightarrow (\exists d \leq a. \text{induced-modulus } d)$
 ⟨*proof*⟩

lemma *conductor-ge-iff*: $\text{conductor} \geq a \longleftrightarrow (\forall d. \text{induced-modulus } d \longrightarrow d \geq a)$
 ⟨*proof*⟩

lemma *conductor-leI*: $\text{induced-modulus } d \implies \text{conductor} \leq d$
 ⟨*proof*⟩

lemma *conductor-geI*: $(\bigwedge d. \text{induced-modulus } d \implies d \geq a) \implies \text{conductor} \geq a$
 ⟨*proof*⟩

lemma *conductor-dvd*: $\text{conductor } \text{dvd } n$

<proof>

lemma *conductor-le-modulus*: $\text{conductor} \leq n$
<proof>

lemma *conductor-gr-0*: $\text{conductor} > 0$
<proof>

lemma *conductor-eq-1-iff-principal*: $\text{conductor} = 1 \iff \chi = \text{principal-dchar } n$
<proof>

lemma *conductor-principal [simp]*: $\chi = \text{principal-dchar } n \implies \text{conductor} = 1$
<proof>

lemma *nonprimitive-imp-conductor-less*:
assumes $\neg \text{primitive-dchar } n \ \chi$
shows $\text{conductor} < n$
<proof>

lemma (in *nonprimitive-dchar*) *conductor-less-modulus*: $\text{conductor} < n$
<proof>

Theorem 8.18

theorem *primitive-principal-form*:
defines $\chi_1 \equiv \text{principal-dchar } n$
assumes $\chi \neq \text{principal-dchar } n$
shows $\exists \Phi. \text{primitive-dchar conductor } \Phi \wedge (\forall n. \chi(n) = \Phi(n) * \chi_1(n))$
<proof>

definition *primitive-extension* :: $\text{nat} \Rightarrow \text{complex}$ **where**
primitive-extension =
(*SOME* $\Phi. \text{primitive-dchar conductor } \Phi \wedge (\forall k. \chi k = \Phi k * \text{principal-dchar } n$
 $k)$)

lemma
assumes *nonprincipal*: $\chi \neq \text{principal-dchar } n$
shows *primitive-primitive-extension*: $\text{primitive-dchar conductor primitive-extension}$
and *principal-decomposition*: $\chi k = \text{primitive-extension } k * \text{principal-dchar } n k$
<proof>

end

7.5 The connection between primitivity and separability

lemma *residue-mult-group-coset*:
fixes $m \ n \ m1 \ m2 :: \text{nat}$ **and** $f :: \text{nat} \Rightarrow \text{nat}$ **and** $G \ H$
defines $G \equiv \text{residue-mult-group } n$
defines $H \equiv \text{residue-mult-group } m$

defines $f \equiv (\lambda k. k \bmod m)$
assumes $b \in (\text{rcosets}_G \text{ kernel } G \ H \ f)$
assumes $m1 \in b \ m2 \in b$
assumes $n > 1 \ m \ \text{dvd} \ n$
shows $m1 \bmod m = m2 \bmod m$
 <proof>

lemma *residue-mult-group-kernel-partition*:
fixes $m \ n :: \text{nat}$ **and** $f :: \text{nat} \Rightarrow \text{nat}$ **and** $G \ H$
defines $G \equiv \text{residue-mult-group } n$
defines $H \equiv \text{residue-mult-group } m$
defines $f \equiv (\lambda k. k \bmod m)$
assumes $m > 1 \ n > 0 \ m \ \text{dvd} \ n$
shows $\text{partition } (\text{carrier } G) (\text{rcosets}_G \text{ kernel } G \ H \ f)$
and $\text{card } (\text{rcosets}_G \text{ kernel } G \ H \ f) = \text{totient } m$
and $\text{card } (\text{kernel } G \ H \ f) = \text{totient } n \ \text{div} \ \text{totient } m$
and $b \in (\text{rcosets}_G \text{ kernel } G \ H \ f) \Longrightarrow b \neq \{\}$
and $b \in (\text{rcosets}_G \text{ kernel } G \ H \ f) \Longrightarrow \text{card } (\text{kernel } G \ H \ f) = \text{card } b$
and $\text{bij-betw } (\lambda b. (\text{the-elem } (f \ ` \ b))) (\text{rcosets}_G \text{ kernel } G \ H \ f) (\text{carrier } H)$
 <proof>

lemma *primitive-iff-separable-lemma*:
assumes $\text{prod}: (\forall n. \chi \ n = \Phi \ n * \chi_1 \ n) \wedge \text{primitive-dchar } d \ \Phi$
assumes $\langle d > 1 \rangle \langle 0 < k \rangle \langle d \ \text{dvd} \ k \rangle \langle k > 1 \rangle$
shows $(\sum m \mid m \in \{1..k\} \wedge \text{coprime } m \ k. \Phi(m) * \text{unity-root } d \ m) =$
 $(\text{totient } k \ \text{div} \ \text{totient } d) * (\sum m \mid m \in \{1..d\} \wedge \text{coprime } m \ d. \Phi(m) * \text{unity-root } d \ m)$
 <proof>

Theorem 8.19

theorem (in *dcharacter*) *primitive-iff-separable*:
 $\text{primitive-dchar } n \ \chi \longleftrightarrow (\forall k > 0. \text{separable } k)$
 <proof>

Theorem 8.20

theorem (in *primitive-dchar*) *fourier-primitive*:
includes *no-vec-lambda-notation*
fixes $\tau :: \text{complex}$
defines $\tau \equiv \text{gauss-sum } 1 / \text{sqrt } n$
shows $\chi \ m = \tau / \text{sqrt } n * (\sum k=1..n. \text{cnj } (\chi \ k) * \text{unity-root } n \ (-m*k))$
and $\text{norm } \tau = 1$
 <proof>

unbundle *vec-lambda-notation*

end

8 The Pólya–Vinogradov Inequality

```
theory Polya-Vinogradov
imports
  Gauss-Sums
  Dirichlet-Series.Divisor-Count
begin

unbundle no-vec-lambda-notation
```

8.1 The case of primitive characters

We first prove a stronger variant of the Pólya–Vinogradov inequality for primitive characters. The fully general variant will then simply be a corollary of this. First, we need some bounds on logarithms, exponentials, and the harmonic numbers:

lemma *ln-add-one-self-less-self*:

```
fixes  $x :: \text{real}$ 
assumes  $x > 0$ 
shows  $\ln (1 + x) < x$ 
⟨proof⟩
```

lemma *exp-1-bounds*:

```
assumes  $x > (0::\text{real})$ 
shows  $\exp 1 > (1 + 1 / x) \text{ powr } x$  and  $\exp 1 < (1 + 1 / x) \text{ powr } (x+1)$ 
⟨proof⟩
```

lemma *harm-aux-ineq-1*:

```
fixes  $k :: \text{real}$ 
assumes  $k > 1$ 
shows  $1 / k < \ln (1 + 1 / (k - 1))$ 
⟨proof⟩
```

lemma *harm-aux-ineq-2-lemma*:

```
assumes  $x \geq (0::\text{real})$ 
shows  $1 < (x + 1) * \ln (1 + 2 / (2 * x + 1))$ 
⟨proof⟩
```

lemma *harm-aux-ineq-2*:

```
fixes  $k :: \text{real}$ 
assumes  $k \geq 1$ 
shows  $1 / (k + 1) < \ln (1 + 2 / (2 * k + 1))$ 
⟨proof⟩
```

lemma *nat-0-1-induct* [*case-names 0 1 step*]:

```
assumes  $P 0 P 1 \wedge n. n \geq 1 \implies P n \implies P (\text{Suc } n)$ 
shows  $P n$ 
⟨proof⟩
```


lemma *harm-less-ln*:
fixes $m :: \text{nat}$
assumes $m > 0$
shows $\text{harm } m < \ln (2 * m + 1)$
 $\langle \text{proof} \rangle$

Theorem 8.21

theorem (*in primitive-dchar*) *polya-vinogradov-inequality-primitive*:
fixes $x :: \text{nat}$
shows $\text{norm } (\sum_{m=1..x} \chi m) < \text{sqrt } n * \ln n$
 $\langle \text{proof} \rangle$

8.2 General case

We now first prove the inequality for the general case in terms of the divisor function:

theorem (*in dcharacter*) *polya-vinogradov-inequality-explicit*:
assumes *nonprincipal*: $\chi \neq \text{principal-dchar } n$
shows $\text{norm } (\text{sum } \chi \{1..x\}) < \text{sqrt conductor} * \ln \text{conductor} * \text{divisor-count } (n \text{ div conductor})$
 $\langle \text{proof} \rangle$

Next, we obtain a suitable upper bound on the number of divisors of n :

lemma *divisor-count-upper-bound-aux*:
fixes $n :: \text{nat}$
shows $\text{divisor-count } n \leq 2 * \text{card } \{d. d \text{ dvd } n \wedge d \leq \text{sqrt } n\}$
 $\langle \text{proof} \rangle$

lemma *divisor-count-upper-bound*:
fixes $n :: \text{nat}$
shows $\text{divisor-count } n \leq 2 * \text{nat } \lfloor \text{sqrt } n \rfloor$
 $\langle \text{proof} \rangle$

lemma *divisor-count-upper-bound'*:
fixes $n :: \text{nat}$
shows $\text{real } (\text{divisor-count } n) \leq 2 * \text{sqrt } n$
 $\langle \text{proof} \rangle$

We are now ready to prove the ‘regular’ Pólya–Vinogradov inequality.

Apostol formulates it in the following way (Theorem 13.15, notation adapted): ‘If χ is any nonprincipal character mod n , then for all $x \geq 2$ we have $\sum_{m \leq x} \chi(m) = O(\sqrt{n} \log n)$.’

The precondition $x \geq 2$ here is completely unnecessary. The ‘Big-O’ notation is somewhat problematic since it does not make explicit in what way the variables are quantified (in particular the x and the χ). The statement of the theorem in this way (for a fixed character χ) seems to suggest that n is

fixed here, which would make the use of ‘Big-O’ completely vacuous, since it is an asymptotic statement about n .

We therefore decided to formulate the inequality in the following more explicit way, even giving an explicit constant factor:

theorem (in d character) *polya-vinogradov-inequality*:
 assumes *nonprincipal: $\chi \neq$ principal- d char n*
 shows *norm $(\sum_{m=1..x} \chi m) < 2 * \text{sqrt } n * \ln n$*
(*proof*)

unbundle *vec-lambda-notation*

end

References

- [1] T. M. Apostol. *Introduction to Analytic Number Theory*. Undergraduate Texts in Mathematics. Springer-Verlag, 1976.