

Free Boolean Algebra

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Abstract

This theory defines a type constructor representing the free Boolean algebra over a set of generators. Values of type $(\alpha)formula$ represent propositional formulas with uninterpreted variables from type α , ordered by implication. In addition to all the standard Boolean algebra operations, the library also provides a function for building homomorphisms to any other Boolean algebra type.

1 Free Boolean algebras

```
theory Free-Boolean-Algebra
imports Main
begin
```

1.1 Free boolean algebra as a set

We start by defining the free boolean algebra over type $'a$ as an inductive set. Here $i :: 'a$ represents a variable; $A :: 'a \text{ set}$ represents a valuation, assigning a truth value to each variable; and $S :: 'a \text{ set set}$ represents a formula, as the set of valuations that make the formula true. The set fb contains representatives of formulas built from finite combinations of variables with negation and conjunction.

inductive-set

```
fb :: 'a set set set
```

where

```
var: {A. i ∈ A} ∈ fb
```

```
| Compl: S ∈ fb ⇒ ¬ S ∈ fb
```

```
| inter: S ∈ fb ⇒ T ∈ fb ⇒ S ∩ T ∈ fb
```

```
lemma fb-Diff: S ∈ fb ⇒ T ∈ fb ⇒ S − T ∈ fb
```

```
unfolding Diff-eq by (intro fb.inter fb.Compl)
```

```
lemma fb-union: S ∈ fb ⇒ T ∈ fb ⇒ S ∪ T ∈ fb
```

```
proof −
```

```
  assume S ∈ fb and T ∈ fb
```

hence $(- S \cap - T) \in fba$ **by** (*intro fba.intros*)
thus $S \cup T \in fba$ **by** *simp*
qed

lemma *fba-empty*: $(\{\} :: 'a \text{ set set}) \in fba$
proof –
obtain $S :: 'a \text{ set set}$ **where** $S \in fba$
by (*fast intro: fba.var*)
hence $S \cap - S \in fba$
by (*intro fba.intros*)
thus *?thesis* **by** *simp*
qed

lemma *fba-UNIV*: $(UNIV :: 'a \text{ set set}) \in fba$
proof –
have $-\{\} \in fba$ **using** *fba-empty* **by** (*rule fba.Compl*)
thus $UNIV \in fba$ **by** *simp*
qed

1.2 Free boolean algebra as a type

The next step is to use *typedef* to define a type isomorphic to the set *fba*. We also define a constructor *var* that corresponds with the similarly-named introduction rule for *fba*.

typedef *'a formula* = *fba* :: *'a set set set*
by (*auto intro: fba-empty*)

definition *var* :: *'a* \Rightarrow *'a formula*
where *var* $i = \text{Abs-formula } \{A. i \in A\}$

lemma *Rep-formula-var*: $\text{Rep-formula } (\text{var } i) = \{A. i \in A\}$
unfolding *var-def* **using** *fba.var* **by** (*rule Abs-formula-inverse*)

Now we make type *'a formula* into a Boolean algebra. This involves defining the various operations (ordering relations, binary infimum and supremum, complement, difference, top and bottom elements) and proving that they satisfy the appropriate laws.

instantiation *formula* :: (*type*) *boolean-algebra*
begin

definition
 $x \sqcap y = \text{Abs-formula } (\text{Rep-formula } x \cap \text{Rep-formula } y)$

definition
 $x \sqcup y = \text{Abs-formula } (\text{Rep-formula } x \cup \text{Rep-formula } y)$

definition
 $\top = \text{Abs-formula } UNIV$

definition

$$\perp = \text{Abs-formula } \{\}$$
definition

$$x \leq y \iff \text{Rep-formula } x \subseteq \text{Rep-formula } y$$
definition

$$x < y \iff \text{Rep-formula } x \subset \text{Rep-formula } y$$
definition

$$- x = \text{Abs-formula } (- \text{Rep-formula } x)$$
definition

$$x - y = \text{Abs-formula } (\text{Rep-formula } x - \text{Rep-formula } y)$$
lemma *Rep-formula-inf*:
$$\text{Rep-formula } (x \sqcap y) = \text{Rep-formula } x \cap \text{Rep-formula } y$$
unfolding *inf-formula-def*
by (*intro Abs-formula-inverse fba.inter Rep-formula*)
lemma *Rep-formula-sup*:
$$\text{Rep-formula } (x \sqcup y) = \text{Rep-formula } x \cup \text{Rep-formula } y$$
unfolding *sup-formula-def*
by (*intro Abs-formula-inverse fba-union Rep-formula*)
lemma *Rep-formula-top*: $\text{Rep-formula } \top = \text{UNIV}$
unfolding *top-formula-def* **by** (*intro Abs-formula-inverse fba-UNIV*)
lemma *Rep-formula-bot*: $\text{Rep-formula } \perp = \{\}$
unfolding *bot-formula-def* **by** (*intro Abs-formula-inverse fba-empty*)
lemma *Rep-formula-compl*: $\text{Rep-formula } (- x) = - \text{Rep-formula } x$ **unfolding** *uminus-formula-def*
by (*intro Abs-formula-inverse fba.Compl Rep-formula*)
lemma *Rep-formula-diff*:
$$\text{Rep-formula } (x - y) = \text{Rep-formula } x - \text{Rep-formula } y$$
unfolding *minus-formula-def*
by (*intro Abs-formula-inverse fba-Diff Rep-formula*)

lemmas *eq-formula-iff* = *Rep-formula-inject* [*symmetric*]
lemmas *Rep-formula-simps* =
less-eq-formula-def less-formula-def eq-formula-iff
Rep-formula-sup Rep-formula-inf Rep-formula-top Rep-formula-bot
Rep-formula-compl Rep-formula-diff Rep-formula-var
instance proof

qed (*unfold Rep-formula-simps, auto*)

end

The laws of a Boolean algebra do not require the top and bottom elements to be distinct, so the following rules must be proved separately:

lemma *bot-neq-top-formula* [*simp*]: $(\perp :: 'a \text{ formula}) \neq \top$
unfolding *Rep-formula-simps* **by** *auto*

lemma *top-neq-bot-formula* [*simp*]: $(\top :: 'a \text{ formula}) \neq \perp$
unfolding *Rep-formula-simps* **by** *auto*

Here we prove an essential property of a free Boolean algebra: all generators are independent.

lemma *var-le-var-simps* [*simp*]:
 $var\ i \leq var\ j \iff i = j$
 $\neg var\ i \leq \neg var\ j$
 $\neg \neg var\ i \leq var\ j$
unfolding *Rep-formula-simps* **by** *fast+*

lemma *var-eq-var-simps* [*simp*]:
 $var\ i = var\ j \iff i = j$
 $var\ i \neq \neg var\ j$
 $\neg var\ i \neq var\ j$
unfolding *Rep-formula-simps set-eq-subset* **by** *fast+*

We conclude this section by proving an induction principle for formulas. It mirrors the definition of the inductive set *fba*, with cases for variables, complements, and conjunction.

lemma *formula-induct* [*case-names var compl inf, induct type: formula*]:
fixes $P :: 'a \text{ formula} \Rightarrow \text{bool}$
assumes 1: $\bigwedge i. P (var\ i)$
assumes 2: $\bigwedge x. P\ x \implies P (\neg x)$
assumes 3: $\bigwedge x\ y. P\ x \implies P\ y \implies P (x \sqcap y)$
shows $P\ x$
proof (*induct x rule: Abs-formula-induct*)
fix $y :: 'a \text{ set set}$
assume $y \in fba$ **thus** $P (Abs-formula\ y)$
proof (*induct rule: fba.induct*)
case (*var i*)
have $P (var\ i)$ **by** (*rule 1*)
thus $?case$ **unfolding** *var-def* .
next
case (*Compl S*)
from $\langle P (Abs-formula\ S) \rangle$ **have** $P (\neg Abs-formula\ S)$ **by** (*rule 2*)
with $\langle S \in fba \rangle$ **show** $?case$
unfolding *uminus-formula-def* **by** (*simp add: Abs-formula-inverse*)

```

next
  case (inter S T)
  from ⟨P (Abs-formula S)⟩ and ⟨P (Abs-formula T)⟩
  have P (Abs-formula S  $\sqcap$  Abs-formula T) by (rule 3)
  with ⟨S  $\in$  fba⟩ and ⟨T  $\in$  fba⟩ show ?case
    unfolding inf-formula-def by (simp add: Abs-formula-inverse)
qed
qed

```

1.3 If-then-else for Boolean algebras

This is a generic if-then-else operator for arbitrary Boolean algebras.

definition

ifte :: 'a::boolean-algebra \Rightarrow 'a \Rightarrow 'a \Rightarrow 'a

where

ifte a x y = (a \sqcap x) \sqcup (\neg a \sqcap y)

lemma ifte-top [simp]: *ifte* \top x y = x

unfolding ifte-def by simp

lemma ifte-bot [simp]: *ifte* \perp x y = y

unfolding ifte-def by simp

lemma ifte-same: *ifte* a x x = x

unfolding ifte-def

by (simp add: inf-sup-distrib2 [symmetric] sup-compl-top)

lemma compl-ifte: \neg *ifte* a x y = *ifte* a (\neg x) (\neg y)

unfolding ifte-def

apply (rule order-antisym)

apply (simp add: inf-sup-distrib1 inf-sup-distrib2 compl-inf-bot)

apply (simp add: sup-inf-distrib1 sup-inf-distrib2 sup-compl-top)

apply (simp add: le-infI1 le-infI2 le-supI1 le-supI2)

apply (simp add: le-infI1 le-infI2 le-supI1 le-supI2)

done

lemma inf-ifte-distrib:

ifte x a b \sqcap *ifte* x c d = *ifte* x (a \sqcap c) (b \sqcap d)

unfolding ifte-def

apply (simp add: inf-sup-distrib1 inf-sup-distrib2)

apply (simp add: inf-sup-aci inf-compl-bot)

done

lemma ifte-ifte-distrib:

ifte x (*ifte* y a b) (*ifte* y c d) = *ifte* y (*ifte* x a c) (*ifte* x b d)

unfolding ifte-def [of x] sup-conv-inf

by (simp only: compl-ifte [symmetric] inf-ifte-distrib [symmetric] ifte-same)

1.4 Formulas over a set of generators

The set *formulas* S consists of those formulas that only depend on variables in the set S . It is analogous to the *lists* operator for the list datatype.

definition

formulas :: 'a set \Rightarrow 'a formula set

where

formulas $S =$
 $\{x. \forall A B. (\forall i \in S. i \in A \longleftrightarrow i \in B) \longrightarrow$
 $A \in \text{Rep-formula } x \longleftrightarrow B \in \text{Rep-formula } x\}$

lemma *formulasI*:

assumes $\bigwedge A B. \forall i \in S. i \in A \longleftrightarrow i \in B$
 $\implies A \in \text{Rep-formula } x \longleftrightarrow B \in \text{Rep-formula } x$

shows $x \in \text{formulas } S$

using *assms unfolding formulas-def by simp*

lemma *formulasD*:

assumes $x \in \text{formulas } S$

assumes $\forall i \in S. i \in A \longleftrightarrow i \in B$

shows $A \in \text{Rep-formula } x \longleftrightarrow B \in \text{Rep-formula } x$

using *assms unfolding formulas-def by simp*

lemma *formulas-mono*: $S \subseteq T \implies \text{formulas } S \subseteq \text{formulas } T$

by (*fast intro!*: *formulasI elim!*: *formulasD*)

lemma *formulas-insert*: $x \in \text{formulas } S \implies x \in \text{formulas } (\text{insert } a \ S)$

unfolding *formulas-def by simp*

lemma *formulas-var*: $i \in S \implies \text{var } i \in \text{formulas } S$

unfolding *formulas-def by (simp add: Rep-formula-simps)*

lemma *formulas-var-iff*: $\text{var } i \in \text{formulas } S \longleftrightarrow i \in S$

unfolding *formulas-def by (simp add: Rep-formula-simps, fast)*

lemma *formulas-bot*: $\perp \in \text{formulas } S$

unfolding *formulas-def by (simp add: Rep-formula-simps)*

lemma *formulas-top*: $\top \in \text{formulas } S$

unfolding *formulas-def by (simp add: Rep-formula-simps)*

lemma *formulas-compl*: $x \in \text{formulas } S \implies \neg x \in \text{formulas } S$

unfolding *formulas-def by (simp add: Rep-formula-simps)*

lemma *formulas-inf*:

$x \in \text{formulas } S \implies y \in \text{formulas } S \implies x \sqcap y \in \text{formulas } S$

unfolding *formulas-def by (auto simp add: Rep-formula-simps)*

lemma *formulas-sup*:

$x \in \text{formulas } S \implies y \in \text{formulas } S \implies x \sqcup y \in \text{formulas } S$
unfolding *formulas-def* **by** (*auto simp add: Rep-formula-simps*)

lemma *formulas-diff*:

$x \in \text{formulas } S \implies y \in \text{formulas } S \implies x - y \in \text{formulas } S$
unfolding *formulas-def* **by** (*auto simp add: Rep-formula-simps*)

lemma *formulas-ifte*:

$a \in \text{formulas } S \implies x \in \text{formulas } S \implies y \in \text{formulas } S \implies$
 $\text{ifte } a \ x \ y \in \text{formulas } S$
unfolding *ifte-def*
by (*intro formulas-sup formulas-inf formulas-compl*)

lemmas *formulas-intros* =

formulas-var formulas-bot formulas-top formulas-compl
formulas-inf formulas-sup formulas-diff formulas-ifte

1.5 Injectivity of if-then-else

The if-then-else operator is injective in some limited circumstances: when the scrutinee is a variable that is not mentioned in either branch.

lemma *ifte-inject*:

assumes $\text{ifte } (\text{var } i) \ x \ y = \text{ifte } (\text{var } i) \ x' \ y'$
assumes $i \notin S$
assumes $x \in \text{formulas } S$ **and** $x' \in \text{formulas } S$
assumes $y \in \text{formulas } S$ **and** $y' \in \text{formulas } S$
shows $x = x' \wedge y = y'$

proof

have 1: $\bigwedge A. i \in A \implies A \in \text{Rep-formula } x \longleftrightarrow A \in \text{Rep-formula } x'$
using *assms(1)*
by (*simp add: Rep-formula-simps ifte-def set-eq-iff, fast*)
have 2: $\bigwedge A. i \notin A \implies A \in \text{Rep-formula } y \longleftrightarrow A \in \text{Rep-formula } y'$
using *assms(1)*
by (*simp add: Rep-formula-simps ifte-def set-eq-iff, fast*)

show $x = x'$

unfolding *Rep-formula-simps*

proof (*rule set-eqI*)

fix A

have $A \in \text{Rep-formula } x \longleftrightarrow \text{insert } i \ A \in \text{Rep-formula } x$

using $\langle x \in \text{formulas } S \rangle$ **by** (*rule formulasD, force simp add: $\langle i \notin S \rangle$*)

also have $\dots \longleftrightarrow \text{insert } i \ A \in \text{Rep-formula } x'$

by (*rule 1, simp*)

also have $\dots \longleftrightarrow A \in \text{Rep-formula } x'$

using $\langle x' \in \text{formulas } S \rangle$ **by** (*rule formulasD, force simp add: $\langle i \notin S \rangle$*)

finally show $A \in \text{Rep-formula } x \longleftrightarrow A \in \text{Rep-formula } x'$.

qed

show $y = y'$

unfolding *Rep-formula-simps*

```

proof (rule set-eqI)
  fix A
  have  $A \in \text{Rep-formula } y \longleftrightarrow A - \{i\} \in \text{Rep-formula } y$ 
    using  $\langle y \in \text{formulas } S \rangle$  by (rule formulasD, force simp add:  $\langle i \notin S \rangle$ )
  also have  $\dots \longleftrightarrow A - \{i\} \in \text{Rep-formula } y'$ 
    by (rule 2, simp)
  also have  $\dots \longleftrightarrow A \in \text{Rep-formula } y'$ 
    using  $\langle y' \in \text{formulas } S \rangle$  by (rule formulasD, force simp add:  $\langle i \notin S \rangle$ )
  finally show  $A \in \text{Rep-formula } y \longleftrightarrow A \in \text{Rep-formula } y'$  .
qed
qed

```

1.6 Specification of homomorphism operator

Our goal is to define a homomorphism operator hom such that for any function f , $hom f$ is the unique Boolean algebra homomorphism satisfying $hom f (var i) = f i$ for all i .

Instead of defining hom directly, we will follow the approach used to define Isabelle's $fold$ operator for finite sets. First we define the graph of the hom function as a relation; later we will define the hom function itself using definite choice.

The hom -graph relation is defined inductively, with introduction rules based on the if-then-else normal form of Boolean formulas. The relation is also indexed by an extra set parameter S , to ensure that branches of each if-then-else do not use the same variable again.

inductive

```

hom-graph ::
  ('a  $\Rightarrow$  'b::boolean-algebra)  $\Rightarrow$  'a set  $\Rightarrow$  'a formula  $\Rightarrow$  'b  $\Rightarrow$  bool
for  $f$  :: 'a  $\Rightarrow$  'b::boolean-algebra

```

where

```

bot: hom-graph  $f$  {} bot bot
| top: hom-graph  $f$  {} top top
| ifte:  $i \notin S \Longrightarrow \text{hom-graph } f S x a \Longrightarrow \text{hom-graph } f S y b \Longrightarrow$ 
  hom-graph  $f$  (insert  $i$   $S$ ) (ifte (var  $i$ )  $x$   $y$ ) (ifte ( $f$   $i$ )  $a$   $b$ )

```

The next two lemmas establish a stronger elimination rule for assumptions of the form hom -graph f (*insert* i S) x a . Essentially, they say that we can arrange the top-level if-then-else to use the variable of our choice. The proof makes use of the distributive properties of if-then-else.

lemma *hom-graph-dest*:

```

hom-graph  $f S x a \Longrightarrow k \in S \Longrightarrow \exists y z b c.$ 
   $x = \text{ifte } (\text{var } k) y z \wedge a = \text{ifte } (f k) b c \wedge$ 
  hom-graph  $f (S - \{k\}) y b \wedge \text{hom-graph } f (S - \{k\}) z c$ 

```

proof (*induct set: hom-graph*)

case (*ifte* i S x a y b) **show** ?*case*

proof (*cases* $i = k$)

assume $i = k$ **with** *ifte*(1,2,4) **show** ?*case* **by** *auto*


```

next
  assume  $i \neq k$ 
  with  $\langle k \in \text{insert } i S \rangle$  have  $k: k \in S$  by simp
  have  $*$ :  $\text{insert } i S - \{k\} = \text{insert } i (S - \{k\})$ 
    using  $\langle i \neq k \rangle$  by (simp add: insert-Diff-if)
  have  $**$ :  $i \notin S - \{k\}$  using  $\langle i \notin S \rangle$  by simp
  from ifte(1) ifte(3) [OF k] ifte(5) [OF k]
  show ?case
    unfolding *
    apply clarify
    apply (simp only: ifte-ifte-distrib [of var i])
    apply (simp only: ifte-ifte-distrib [of f i])
    apply (fast intro: hom-graph.ifte [OF **])
    done
qed
qed simp-all

```

```

lemma hom-graph-insert-elim:
  assumes hom-graph f (insert i S) x a and  $i \notin S$ 
  obtains  $y z b c$ 
  where  $x = \text{ifte } (\text{var } i) y z$ 
    and  $a = \text{ifte } (f i) b c$ 
    and hom-graph f S y b
    and hom-graph f S z c
using hom-graph-dest [OF assms(1) insertI1]
by (clarify, simp add: assms(2))

```

Now we prove the first uniqueness property of the *hom-graph* relation. This version of uniqueness says that for any particular value of S , the relation $\text{hom-graph } f S$ maps each x to at most one a . The proof uses the injectiveness of if-then-else, which we proved earlier.

```

lemma hom-graph-imp-formulas:
  hom-graph f S x a  $\implies x \in \text{formulas } S$ 
by (induct set: hom-graph, simp-all add: formulas-intros formulas-insert)

```

```

lemma hom-graph-unique:
  hom-graph f S x a  $\implies \text{hom-graph } f S x a' \implies a = a'$ 
proof (induct arbitrary:  $a'$  set: hom-graph)
  case (ifte i S y b z c a')
  from ifte(6,1) obtain  $y' z' b' c'$ 
  where 1:  $\text{ifte } (\text{var } i) y z = \text{ifte } (\text{var } i) y' z'$ 
    and 2:  $a' = \text{ifte } (f i) b' c'$ 
    and 3: hom-graph f S y' b'
    and 4: hom-graph f S z' c'
  by (rule hom-graph-insert-elim)
  from 1 3 4 ifte(1,2,4) have  $y = y' \wedge z = z'$ 
  by (intro ifte-inject hom-graph-imp-formulas)
  with 2 3 4 ifte(3,5) show ifte (f i) b c = a'
  by simp

```

qed (*erule hom-graph.cases, simp-all*)⁺

The next few lemmas will help to establish a stronger version of the uniqueness property of *hom-graph*. They show that the *hom-graph* relation is preserved if we replace S with a larger finite set.

lemma *hom-graph-insert*:
assumes *hom-graph f S x a*
shows *hom-graph f (insert i S) x a*
proof (*cases i ∈ S*)
assume $i \in S$ **with** *assms* **show** *?thesis* **by** (*simp add: insert-absorb*)
next
assume $i \notin S$
hence *hom-graph f (insert i S) (ifte (var i) x x) (ifte (f i) a a)*
by (*intro hom-graph.ifte assms*)
thus *hom-graph f (insert i S) x a*
by (*simp only: ifte-same*)
qed

lemma *hom-graph-finite-superset*:
assumes *hom-graph f S x a* **and** *finite T* **and** $S \subseteq T$
shows *hom-graph f T x a*
proof –
from $\langle \text{finite } T \rangle$ **have** *hom-graph f (S ∪ T) x a*
by (*induct set: finite, simp add: assms, simp add: hom-graph-insert*)
with $\langle S \subseteq T \rangle$ **show** *hom-graph f T x a*
by (*simp only: subset-Un-eq*)
qed

lemma *hom-graph-imp-finite*:
hom-graph f S x a \implies *finite S*
by (*induct set: hom-graph*) *simp-all*

This stronger uniqueness property says that *hom-graph f* maps each x to at most one a , even for *different* values of the set parameter.

lemma *hom-graph-unique'*:
assumes *hom-graph f S x a* **and** *hom-graph f T x a'*
shows $a = a'$
proof (*rule hom-graph-unique*)
have *fin: finite (S ∪ T)*
using *assms* **by** (*intro finite-UnI hom-graph-imp-finite*)
show *hom-graph f (S ∪ T) x a*
using *assms(1) fin Un-upper1* **by** (*rule hom-graph-finite-superset*)
show *hom-graph f (S ∪ T) x a'*
using *assms(2) fin Un-upper2* **by** (*rule hom-graph-finite-superset*)
qed

Finally, these last few lemmas establish that the *hom-graph f* relation is total: every x is mapped to some a .

```

lemma hom-graph-var: hom-graph  $f$   $\{i\}$  (var  $i$ ) ( $f$   $i$ )
proof –
  have hom-graph  $f$   $\{i\}$  (ifte (var  $i$ ) top bot) (ifte ( $f$   $i$ ) top bot)
    by (simp add: hom-graph.intros)
  thus hom-graph  $f$   $\{i\}$  (var  $i$ ) ( $f$   $i$ )
    unfolding ifte-def by simp
qed

lemma hom-graph-compl:
  hom-graph  $f$   $S$   $x$   $a$   $\implies$  hom-graph  $f$   $S$  ( $-$   $x$ ) ( $-$   $a$ )
by (induct set: hom-graph, simp-all add: hom-graph.intros compl-ifte)

lemma hom-graph-inf:
  hom-graph  $f$   $S$   $x$   $a$   $\implies$  hom-graph  $f$   $S$   $y$   $b$   $\implies$ 
  hom-graph  $f$   $S$  ( $x$   $\sqcap$   $y$ ) ( $a$   $\sqcap$   $b$ )
apply (induct arbitrary:  $y$   $b$  set: hom-graph)
apply (simp add: hom-graph.bot)
apply simp
apply (erule (1) hom-graph-insert-elim)
apply (auto simp add: inf-ifte-distrib hom-graph.ifte)
done

lemma hom-graph-union-inf:
  assumes hom-graph  $f$   $S$   $x$   $a$  and hom-graph  $f$   $T$   $y$   $b$ 
  shows hom-graph  $f$  ( $S$   $\cup$   $T$ ) ( $x$   $\sqcap$   $y$ ) ( $a$   $\sqcap$   $b$ )
proof (rule hom-graph-inf)
  have fin: finite ( $S$   $\cup$   $T$ )
    using assms by (intro finite-UnI hom-graph-imp-finite)
  show hom-graph  $f$  ( $S$   $\cup$   $T$ )  $x$   $a$ 
    using assms(1) fin Un-upper1 by (rule hom-graph-finite-superset)
  show hom-graph  $f$  ( $S$   $\cup$   $T$ )  $y$   $b$ 
    using assms(2) fin Un-upper2 by (rule hom-graph-finite-superset)
qed

lemma hom-graph-exists:  $\exists a$   $S$ . hom-graph  $f$   $S$   $x$   $a$ 
by (induct  $x$ )
  (auto intro: hom-graph-var hom-graph-compl hom-graph-union-inf)

```

1.7 Homomorphisms into other boolean algebras

Now that we have proved the necessary existence and uniqueness properties of *hom-graph*, we can define the function *hom* using definite choice.

definition

hom :: ($'a \Rightarrow 'b$:*boolean-algebra*) \Rightarrow $'a$ *formula* \Rightarrow $'b$

where

hom f x = (*THE* a . $\exists S$. *hom-graph* f S x a)

lemma *hom-graph-hom*: $\exists S$. *hom-graph* f S x (*hom* f x)

unfolding *hom-def*

```

apply (rule theI')
apply (rule ex-ex1I)
apply (rule hom-graph-exists)
apply (fast elim: hom-graph-unique')
done

```

```

lemma hom-equality:
  hom-graph f S x a  $\implies$  hom f x = a
unfolding hom-def
apply (rule the-equality)
apply (erule exI)
apply (erule exE)
apply (erule (1) hom-graph-unique')
done

```

The *hom* function correctly implements its specification:

```

lemma hom-var [simp]: hom f (var i) = f i
by (rule hom-equality, rule hom-graph-var)

```

```

lemma hom-bot [simp]: hom f  $\perp$  =  $\perp$ 
by (rule hom-equality, rule hom-graph.bot)

```

```

lemma hom-top [simp]: hom f  $\top$  =  $\top$ 
by (rule hom-equality, rule hom-graph.top)

```

```

lemma hom-compl [simp]: hom f ( $-$  x) =  $-$  hom f x

```

```

proof –
  obtain S where hom-graph f S x (hom f x)
  using hom-graph-hom ..
  hence hom-graph f S ( $-$  x) ( $-$  hom f x)
  by (rule hom-graph-compl)
  thus hom f ( $-$  x) =  $-$  hom f x
  by (rule hom-equality)
qed

```

```

lemma hom-inf [simp]: hom f (x  $\sqcap$  y) = hom f x  $\sqcap$  hom f y

```

```

proof –
  obtain S where S: hom-graph f S x (hom f x)
  using hom-graph-hom ..
  obtain T where T: hom-graph f T y (hom f y)
  using hom-graph-hom ..
  have hom-graph f (S  $\cup$  T) (x  $\sqcap$  y) (hom f x  $\sqcap$  hom f y)
  using S T by (rule hom-graph-union-inf)
  thus ?thesis by (rule hom-equality)
qed

```

```

lemma hom-sup [simp]: hom f (x  $\sqcup$  y) = hom f x  $\sqcup$  hom f y
unfolding sup-conv-inf by (simp only: hom-compl hom-inf)

```

lemma *hom-diff* [*simp*]: $\text{hom } f (x - y) = \text{hom } f x - \text{hom } f y$
unfolding *diff-eq* **by** (*simp only: hom-compl hom-inf*)

lemma *hom-ifte* [*simp*]:
 $\text{hom } f (\text{ifte } x y z) = \text{ifte } (\text{hom } f x) (\text{hom } f y) (\text{hom } f z)$
unfolding *ifte-def* **by** (*simp only: hom-compl hom-inf hom-sup*)

lemmas *hom-simps* =
hom-var hom-bot hom-top hom-compl
hom-inf hom-sup hom-diff hom-ifte

The type *'a formula* can be viewed as a monad, with *var* as the unit, and *hom* as the bind operator. We can prove the standard monad laws with simple proofs by induction.

lemma *hom-var-eq-id*: $\text{hom } \text{var } x = x$
by (*induct x*) *simp-all*

lemma *hom-hom*: $\text{hom } f (\text{hom } g x) = \text{hom } (\lambda i. \text{hom } f (g i)) x$
by (*induct x*) *simp-all*

1.8 Map operation on Boolean formulas

We can define a map functional in terms of *hom* and *var*. The properties of *fmap* follow directly from the lemmas we have already proved about *hom*.

definition
 $\text{fmap} :: ('a \Rightarrow 'b) \Rightarrow 'a \text{ formula} \Rightarrow 'b \text{ formula}$
where
 $\text{fmap } f = \text{hom } (\lambda i. \text{var } (f i))$

lemma *fmap-var* [*simp*]: $\text{fmap } f (\text{var } i) = \text{var } (f i)$
unfolding *fmap-def* **by** *simp*

lemma *fmap-bot* [*simp*]: $\text{fmap } f \perp = \perp$
unfolding *fmap-def* **by** *simp*

lemma *fmap-top* [*simp*]: $\text{fmap } f \top = \top$
unfolding *fmap-def* **by** *simp*

lemma *fmap-compl* [*simp*]: $\text{fmap } f (- x) = - \text{fmap } f x$
unfolding *fmap-def* **by** *simp*

lemma *fmap-inf* [*simp*]: $\text{fmap } f (x \sqcap y) = \text{fmap } f x \sqcap \text{fmap } f y$
unfolding *fmap-def* **by** *simp*

lemma *fmap-sup* [*simp*]: $\text{fmap } f (x \sqcup y) = \text{fmap } f x \sqcup \text{fmap } f y$
unfolding *fmap-def* **by** *simp*

lemma *fmap-diff* [*simp*]: $\text{fmap } f (x - y) = \text{fmap } f x - \text{fmap } f y$

unfolding *fmap-def* **by** *simp*

lemma *fmap-ifte* [*simp*]:

$fmap\ f\ (ifte\ x\ y\ z) = ifte\ (fmap\ f\ x)\ (fmap\ f\ y)\ (fmap\ f\ z)$

unfolding *fmap-def* **by** *simp*

lemmas *fmap-simps* =

fmap-var fmap-bot fmap-top fmap-compl

fmap-inf fmap-sup fmap-diff fmap-ifte

The map functional satisfies the functor laws: it preserves identity and function composition.

lemma *fmap-ident*: $fmap\ (\lambda i. i)\ x = x$

by (*induct x*) *simp-all*

lemma *fmap-fmap*: $fmap\ f\ (fmap\ g\ x) = fmap\ (f \circ g)\ x$

by (*induct x*) *simp-all*

1.9 Hiding lattice syntax

The following command hides the lattice syntax, to avoid potential conflicts with other theories that import this one. To re-enable the syntax, users should unbundle *lattice-syntax*.

unbundle *no-lattice-syntax*

end