

Finite Fields

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June 16, 2022

Abstract

This entry formalizes the classification of the finite fields (also called Galois fields): For each prime power p^n there exists exactly one (up to isomorphisms) finite field of that size and there are no other finite fields. The derivation includes a formalization of the characteristic of rings, the Frobenius endomorphism, formal differentiation for polynomials in HOL-Algebra and Gauss' formula for the number of monic irreducible polynomials over finite fields:

$$\frac{1}{n} \sum_{d|n} \mu(d) p^{n/d}.$$

The proofs are based on the books from Ireland and Rosen [3], as well as, Lidl and Niederreiter [4].

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1 Introduction

The following section starts with preliminary results. Section 3 introduces the characteristic of rings with the Frobenius endomorphism. Whenever it makes sense, the definitions and facts do not assume the finiteness of the fields or rings. For example the characteristic is defined over arbitrary rings (and also fields). While formal derivatives do exist for type-class based structures in `HOL-Computational_Algebra`, as far as I can tell, they do not exist for the structure based polynomials in `HOL-Algebra`. These are introduced in Section 4.

A cornerstone of the proof is the derivation of Gauss' formula for the number of monic irreducible polynomials over a finite field R in Section 6.2. The proof follows the derivation by Ireland and Rosen [3, §7] closely, with the caveat that it does not assume that R is a simple prime field, but that it is just a finite field. This works by adjusting a proof step with the information that the order of a finite field must be of the form p^n , where p is the characteristic of the field, derived in Section 3. The final step relies on the Möbius inversion theorem formalized by Eberl [2].¹ With Gauss' formula it is possible to show the existence of the finite fields of order p^n where p is a prime and $n > 0$. During the proof the fact that the polynomial $X^n - X$ splits in a field of order n is also derived, which is necessary for the uniqueness result as well.

The uniqueness proof is inspired by the derivation of the same result in Lidl and Niederreiter [4], but because of the already derived existence proof for irreducible polynomials, it was possible to reduce its complexity.

The classification consists of three theorems:

- *Existence*: For each prime power p^n there exists a finite field of that size. This is shown at the conclusion of Section 6.2.
- *Uniqueness*: Any two finite fields of the same size are isomorphic. This is shown at the conclusion of Section 7.
- *Completeness*: Any finite fields' size must be a prime power. This is shown at the conclusion of Section 3.

¹Thanks to Katharina Kreuzer for discovering that formalization.

2 Preliminary Results

theory *Finite-Fields-Preliminary-Results*
imports *HOL-Algebra.Polynomial-Divisibility*
begin

2.1 Summation in the discrete topology

The following lemmas transfer the corresponding result from the summation over finite sets to summation over functions which vanish outside of a finite set.

lemma *sum'-subtractf-nat*:
fixes $f :: 'a \Rightarrow \text{nat}$
assumes $\text{finite } \{i \in A. f\ i \neq 0\}$
assumes $\bigwedge i. i \in A \implies g\ i \leq f\ i$
shows $\text{sum}' (\lambda i. f\ i - g\ i)\ A = \text{sum}' f\ A - \text{sum}' g\ A$
(**is** *?lhs = ?rhs*)
<proof>

lemma *sum'-nat-eq-0-iff*:
fixes $f :: 'a \Rightarrow \text{nat}$
assumes $\text{finite } \{i \in A. f\ i \neq 0\}$
assumes $\text{sum}' f\ A = 0$
shows $\bigwedge i. i \in A \implies f\ i = 0$
<proof>

lemma *sum'-eq-iff*:
fixes $f :: 'a \Rightarrow \text{nat}$
assumes $\text{finite } \{i \in A. f\ i \neq 0\}$
assumes $\bigwedge i. i \in A \implies f\ i \geq g\ i$
assumes $\text{sum}' f\ A \leq \text{sum}' g\ A$
shows $\forall i \in A. f\ i = g\ i$
<proof>

2.2 Polynomials

The embedding of the constant polynomials into the polynomials is injective:

lemma (**in** *ring*) *poly-of-const-inj*:
inj poly-of-const
<proof>

lemma (**in** *domain*) *embed-hom*:
assumes *subring* $K\ R$
shows *ring-hom-ring* $(K[X])\ (\text{poly-ring } R)\ \text{id}$
<proof>

The following are versions of the properties of the degrees of poly-

nomials, that abstract over the definition of the polynomial ring structure. In the theories *HOL–Algebra.Polynomials* and also *HOL–Algebra.Polynomial-Divisibility* these abstract version are usually indicated with the suffix “shell”, consider for example: *domain.pdivides-iff-shell*.

lemma (in ring) *degree-add-distinct*:
assumes *subring* $K R$
assumes $f \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
assumes $g \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
assumes $\text{degree } f \neq \text{degree } g$
shows $\text{degree } (f \oplus_{K[X]} g) = \max (\text{degree } f) (\text{degree } g)$
 $\langle \text{proof} \rangle$

lemma (in domain) *degree-mult*:
assumes *subring* $K R$
assumes $f \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
assumes $g \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
shows $\text{degree } (f \otimes_{K[X]} g) = \text{degree } f + \text{degree } g$
 $\langle \text{proof} \rangle$

lemma (in ring) *degree-one*:
 $\text{degree } (\mathbf{1}_{K[X]}) = 0$
 $\langle \text{proof} \rangle$

lemma (in domain) *pow-non-zero*:
 $x \in \text{carrier } R \implies x \neq \mathbf{0} \implies x [\wedge] (n :: \text{nat}) \neq \mathbf{0}$
 $\langle \text{proof} \rangle$

lemma (in domain) *degree-pow*:
assumes *subring* $K R$
assumes $f \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
shows $\text{degree } (f [\wedge]_{K[X]} n) = \text{degree } f * n$
 $\langle \text{proof} \rangle$

lemma (in ring) *degree-var*:
 $\text{degree } (X_R) = 1$
 $\langle \text{proof} \rangle$

lemma (in domain) *var-carr*:
fixes $n :: \text{nat}$
assumes *subring* $K R$
shows $X_R \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
 $\langle \text{proof} \rangle$

lemma (in domain) *var-pow-carr*:
fixes $n :: \text{nat}$
assumes *subring* $K R$

shows $X_R [\wedge]_K [X] n \in \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
 <proof>

lemma (in domain) *var-pow-degree*:
fixes $n :: \text{nat}$
assumes *subring* $K R$
shows $\text{degree } (X_R [\wedge]_K [X] n) = n$
 <proof>

lemma (in domain) *finprod-non-zero*:
assumes *finite* A
assumes $f \in A \rightarrow \text{carrier } R - \{\mathbf{0}\}$
shows $(\otimes_{i \in A} f i) \in \text{carrier } R - \{\mathbf{0}\}$
 <proof>

lemma (in domain) *degree-prod*:
assumes *finite* A
assumes *subring* $K R$
assumes $f \in A \rightarrow \text{carrier } (K[X]) - \{\mathbf{0}_{K[X]}\}$
shows $\text{degree } (\otimes_{K[X] i \in A} f i) = (\sum_{i \in A} \text{degree } (f i))$
 <proof>

lemma (in ring) *coeff-add*:
assumes *subring* $K R$
assumes $f \in \text{carrier } (K[X]) g \in \text{carrier } (K[X])$
shows $\text{coeff } (f \oplus_{K[X]} g) i = \text{coeff } f i \oplus_R \text{coeff } g i$
 <proof>

This is a version of geometric sums for commutative rings:

lemma (in cring) *geom*:
fixes $q :: \text{nat}$
assumes [*simp*]: $a \in \text{carrier } R$
shows $(a \ominus \mathbf{1}) \otimes (\bigoplus_{i \in \{..<q\}} a [\wedge] i) = (a [\wedge] q \ominus \mathbf{1})$
 (is ?lhs = ?rhs)
 <proof>

lemma (in domain) *rupture-eq-0-iff*:
assumes *subfield* $K R$ $p \in \text{carrier } (K[X]) q \in \text{carrier } (K[X])$
shows $\text{rupture-surj } K p q = \mathbf{0}_{\text{Rupt } K p} \longleftrightarrow p \text{ pdivides } q$
 (is ?lhs \longleftrightarrow ?rhs)
 <proof>

2.3 Ring Isomorphisms

The following lemma shows that an isomorphism between domains also induces an isomorphism between the corresponding polynomial rings.

lemma *lift-iso-to-poly-ring*:

assumes $h \in \text{ring-iso } R \ S \ \text{domain } R \ \text{domain } S$
shows $\text{map } h \in \text{ring-iso } (\text{poly-ring } R) \ (\text{poly-ring } S)$
 $\langle \text{proof} \rangle$

lemma *carrier-hom*:
assumes $f \in \text{carrier } (\text{poly-ring } R)$
assumes $h \in \text{ring-iso } R \ S \ \text{domain } R \ \text{domain } S$
shows $\text{map } h \ f \in \text{carrier } (\text{poly-ring } S)$
 $\langle \text{proof} \rangle$

lemma *carrier-hom'*:
assumes $f \in \text{carrier } (\text{poly-ring } R)$
assumes $h \in \text{ring-hom } R \ S$
assumes $\text{domain } R \ \text{domain } S$
assumes $\text{inj-on } h \ (\text{carrier } R)$
shows $\text{map } h \ f \in \text{carrier } (\text{poly-ring } S)$
 $\langle \text{proof} \rangle$

The following lemmas transfer properties like divisibility, irreducibility etc. between ring isomorphisms.

lemma *divides-hom*:
assumes $h \in \text{ring-iso } R \ S$
assumes $\text{domain } R \ \text{domain } S$
assumes $x \in \text{carrier } R \ y \in \text{carrier } R$
shows $x \ \text{divides}_R \ y \iff (h \ x) \ \text{divides}_S \ (h \ y) \ (\text{is } ?lhs \iff ?rhs)$
 $\langle \text{proof} \rangle$

lemma *properfactor-hom*:
assumes $h \in \text{ring-iso } R \ S$
assumes $\text{domain } R \ \text{domain } S$
assumes $x \in \text{carrier } R \ b \in \text{carrier } R$
shows $\text{properfactor } R \ b \ x \iff \text{properfactor } S \ (h \ b) \ (h \ x)$
 $\langle \text{proof} \rangle$

lemma *Units-hom*:
assumes $h \in \text{ring-iso } R \ S$
assumes $\text{domain } R \ \text{domain } S$
assumes $x \in \text{carrier } R$
shows $x \in \text{Units } R \iff h \ x \in \text{Units } S$
 $\langle \text{proof} \rangle$

lemma *irreducible-hom*:
assumes $h \in \text{ring-iso } R \ S$
assumes $\text{domain } R \ \text{domain } S$
assumes $x \in \text{carrier } R$
shows $\text{irreducible } R \ x = \text{irreducible } S \ (h \ x)$
 $\langle \text{proof} \rangle$

lemma *pirreducible-hom*:

assumes $h \in \text{ring-iso } R \ S$
assumes $\text{domain } R \ \text{domain } S$
assumes $f \in \text{carrier } (\text{poly-ring } R)$
shows $\text{pirreducible}_R (\text{carrier } R) f =$
 $\text{pirreducible}_S (\text{carrier } S) (\text{map } h f)$
(is $?lhs = ?rhs)$
 $\langle \text{proof} \rangle$

lemma *ring-hom-cong*:
assumes $\bigwedge x. x \in \text{carrier } R \implies f' x = f x$
assumes *ring* R
assumes $f \in \text{ring-hom } R \ S$
shows $f' \in \text{ring-hom } R \ S$
 $\langle \text{proof} \rangle$

The natural homomorphism between factor rings, where one ideal is a subset of the other.

lemma (*in ring*) *quot-quot-hom*:
assumes *ideal* $I \ R$
assumes *ideal* $J \ R$
assumes $I \subseteq J$
shows $(\lambda x. (J \langle + \rangle_R x)) \in \text{ring-hom } (R \ \text{Quot } I) \ (R \ \text{Quot } J)$
 $\langle \text{proof} \rangle$

lemma (*in ring*) *quot-carr*:
assumes *ideal* $I \ R$
assumes $y \in \text{carrier } (R \ \text{Quot } I)$
shows $y \subseteq \text{carrier } R$
 $\langle \text{proof} \rangle$

lemma (*in ring*) *set-add-zero*:
assumes $A \subseteq \text{carrier } R$
shows $\{0\} \langle + \rangle_R A = A$
 $\langle \text{proof} \rangle$

Adapted from the proof of *domain.polynomial-rupture*

lemma (*in domain*) *rupture-surj-as-eval*:
assumes *subring* $K \ R$
assumes $p \in \text{carrier } (K[X]) \ q \in \text{carrier } (K[X])$
shows $\text{rupture-surj } K \ p \ q =$
 $\text{ring.eval } (\text{Rupt } K \ p) (\text{map } ((\text{rupture-surj } K \ p) \circ \text{poly-of-const}) \ q)$
 $(\text{rupture-surj } K \ p \ X)$
 $\langle \text{proof} \rangle$

2.4 Divisibility

lemma (*in field*) *f-comm-group-1*:
assumes $x \in \text{carrier } R \ y \in \text{carrier } R$

assumes $x \neq \mathbf{0} \ y \neq \mathbf{0}$
assumes $x \otimes y = \mathbf{0}$
shows *False*
 <proof>

lemma (in *field*) *f-comm-group-2*:
assumes $x \in \text{carrier } R$
assumes $x \neq \mathbf{0}$
shows $\exists y \in \text{carrier } R - \{\mathbf{0}\}. y \otimes x = \mathbf{1}$
 <proof>

sublocale *field < mult-of: comm-group mult-of R*
rewrites $\text{mult } (\text{mult-of } R) = \text{mult } R$
and one $(\text{mult-of } R) = \text{one } R$
 <proof>

lemma (in *domain*) *div-neg*:
assumes $a \in \text{carrier } R \ b \in \text{carrier } R$
assumes $a \text{ divides } b$
shows $a \text{ divides } (\ominus b)$
 <proof>

lemma (in *domain*) *div-sum*:
assumes $a \in \text{carrier } R \ b \in \text{carrier } R \ c \in \text{carrier } R$
assumes $a \text{ divides } b$
assumes $a \text{ divides } c$
shows $a \text{ divides } (b \oplus c)$
 <proof>

lemma (in *domain*) *div-sum-iff*:
assumes $a \in \text{carrier } R \ b \in \text{carrier } R \ c \in \text{carrier } R$
assumes $a \text{ divides } b$
shows $a \text{ divides } (b \oplus c) \longleftrightarrow a \text{ divides } c$
 <proof>

end

2.5 Factorization

theory *Finite-Fields-Factorization-Ext*
imports *Finite-Fields-Preliminary-Results*
begin

This section contains additional results building on top of the development in *HOL-Algebra.Divisibility* about factorization in a *factorial-monoid*.

definition *factor-mset* **where** *factor-mset* $G \ x =$
 (*THE* $f. (\exists \text{ as. } f = \text{fmset } G \ \text{as} \wedge \text{wfactors } G \ \text{as } x \wedge \text{set as} \subseteq \text{carrier } G)$)

In *HOL–Algebra.Divisibility* it is already verified that the multiset representing the factorization of an element of a factorial monoid into irreducible factors is well-defined. With these results it is then possible to define *factor-mset* and show its properties, without referring to a factorization in list form first.

definition *multiplicity where*

multiplicity G d $g = \text{Max } \{(n::\text{nat}). (d [\wedge]_G n) \text{ divides}_G g\}$

definition *canonical-irreducibles where*

canonical-irreducibles G $A = ($
 $A \subseteq \{a. a \in \text{carrier } G \wedge \text{irreducible } G a\} \wedge$
 $(\forall x y. x \in A \longrightarrow y \in A \longrightarrow x \sim_G y \longrightarrow x = y) \wedge$
 $(\forall x \in \text{carrier } G. \text{irreducible } G x \longrightarrow (\exists y \in A. x \sim_G y)))$

A set of irreducible elements that contains exactly one element from each equivalence class of an irreducible element formed by association, is called a set of *canonical-irreducibles*. An example is the set of monic irreducible polynomials as representatives of all irreducible polynomials.

context *factorial-monoid*

begin

lemma *assoc-as-fmset-eq:*

assumes *wfactors* G *as* a
and *wfactors* G *bs* b
and $a \in \text{carrier } G$
and $b \in \text{carrier } G$
and $\text{set } as \subseteq \text{carrier } G$
and $\text{set } bs \subseteq \text{carrier } G$
shows $a \sim b \longleftrightarrow (\text{fmset } G as = \text{fmset } G bs)$

<proof>

lemma *factor-mset-aux-1:*

assumes $a \in \text{carrier } G$ $\text{set } as \subseteq \text{carrier } G$ *wfactors* G *as* a
shows $\text{factor-mset } G a = \text{fmset } G as$

<proof>

lemma *factor-mset-aux:*

assumes $a \in \text{carrier } G$
shows $\exists as. \text{factor-mset } G a = \text{fmset } G as \wedge \text{wfactors } G as a \wedge$
 $\text{set } as \subseteq \text{carrier } G$

<proof>

lemma *factor-mset-set:*

assumes $a \in \text{carrier } G$
assumes $x \in \# \text{factor-mset } G a$
obtains y **where**
 $y \in \text{carrier } G$

irreducible G y
assocs G y = x
 ⟨proof⟩

lemma *factor-mset-mult*:
assumes $a \in \text{carrier } G$ $b \in \text{carrier } G$
shows $\text{factor-mset } G (a \otimes b) = \text{factor-mset } G a + \text{factor-mset } G b$
 ⟨proof⟩

lemma *factor-mset-unit*: $\text{factor-mset } G \mathbf{1} = \{\#\}$
 ⟨proof⟩

lemma *factor-mset-irred*:
assumes $x \in \text{carrier } G$ *irreducible G x*
shows $\text{factor-mset } G x = \text{image-mset } (\text{assocs } G) \{\#x\#\}$
 ⟨proof⟩

lemma *factor-mset-divides*:
assumes $a \in \text{carrier } G$ $b \in \text{carrier } G$
shows $a \text{ divides } b \iff \text{factor-mset } G a \subseteq\# \text{factor-mset } G b$
 ⟨proof⟩

lemma *factor-mset-sim*:
assumes $a \in \text{carrier } G$ $b \in \text{carrier } G$
shows $a \sim b \iff \text{factor-mset } G a = \text{factor-mset } G b$
 ⟨proof⟩

lemma *factor-mset-prod*:
assumes *finite A*
assumes $f ' A \subseteq \text{carrier } G$
shows $\text{factor-mset } G (\bigotimes a \in A. f a) =$
 $(\sum a \in A. \text{factor-mset } G (f a))$
 ⟨proof⟩

lemma *factor-mset-pow*:
assumes $a \in \text{carrier } G$
shows $\text{factor-mset } G (a [\wedge] n) = \text{repeat-mset } n (\text{factor-mset } G a)$
 ⟨proof⟩

lemma *image-mset-sum*:
assumes *finite F*
shows
 $\text{image-mset } h (\sum x \in F. f x) = (\sum x \in F. \text{image-mset } h (f x))$
 ⟨proof⟩

lemma *decomp-mset*:
 $(\sum x \in \text{set-mset } R. \text{replicate-mset } (\text{count } R x) x) = R$
 ⟨proof⟩

lemma *factor-mset-count*:

assumes $a \in \text{carrier } G$ $d \in \text{carrier } G$ *irreducible* G d

shows $\text{count } (\text{factor-mset } G a) (\text{assocs } G d) = \text{multiplicity } G d a$

<proof>

lemma *multiplicity-ge-iff*:

assumes $d \in \text{carrier } G$ *irreducible* G d $a \in \text{carrier } G$

shows $\text{multiplicity } G d a \geq k \iff d \text{ [}\hat{\text{]}} k \text{ divides } a$

(**is** $?lhs \iff ?rhs$)

<proof>

lemma *multiplicity-gt-0-iff*:

assumes $d \in \text{carrier } G$ *irreducible* G d $a \in \text{carrier } G$

shows $\text{multiplicity } G d a > 0 \iff d \text{ divides } a$

<proof>

lemma *factor-mset-count-2*:

assumes $a \in \text{carrier } G$

assumes $\bigwedge z. z \in \text{carrier } G \implies \text{irreducible } G z \implies y \neq \text{assocs } G z$

shows $\text{count } (\text{factor-mset } G a) y = 0$

<proof>

lemma *factor-mset-choose*:

assumes $a \in \text{carrier } G$ *set-mset* $R \subseteq \text{carrier } G$

assumes $\text{image-mset } (\text{assocs } G) R = \text{factor-mset } G a$

shows $a \sim (\bigotimes_{x \in \text{set-mset } R} x \text{ [}\hat{\text{]}} \text{count } R x)$ (**is** $a \sim ?rhs$)

<proof>

lemma *divides-iff-mult-mono*:

assumes $a \in \text{carrier } G$ $b \in \text{carrier } G$

assumes *canonical-irreducibles* G R

assumes $\bigwedge d. d \in R \implies \text{multiplicity } G d a \leq \text{multiplicity } G d b$

shows $a \text{ divides } b$

<proof>

lemma *count-image-mset-inj*:

assumes *inj-on* f R $x \in R$ *set-mset* $A \subseteq R$

shows $\text{count } (\text{image-mset } f A) (f x) = \text{count } A x$

<proof>

Factorization of an element from a *factorial-monoid* using a selection of representatives from each equivalence class formed by (\sim) .

lemma *split-factors*:

assumes *canonical-irreducibles* G R

assumes $a \in \text{carrier } G$

shows

finite $\{d. d \in R \wedge \text{multiplicity } G d a > 0\}$

$a \sim (\bigotimes_{d \in \{d. d \in R \wedge \text{multiplicity } G d a > 0\}} d)$.

```

    d [^] multiplicity G d a) (is a ~ ?rhs)
  <proof>

end

end

```

3 Characteristic of Rings

```

theory Ring-Characteristic
  imports
    Finite-Fields-Factorization-Ext
    HOL-Algebra.IntRing
    HOL-Algebra.Embedded-Algebras
begin

  locale finite-field = field +
    assumes finite-carrier: finite (carrier R)
  begin

    lemma finite-field-min-order:
      order R > 1
    <proof>

    lemma (in finite-field) order-pow-eq-self:
      assumes x ∈ carrier R
      shows x [^] (order R) = x
    <proof>

    lemma (in finite-field) order-pow-eq-self':
      assumes x ∈ carrier R
      shows x [^] (order R ^ d) = x
    <proof>

  end

  lemma finite-fieldI:
    assumes field R
    assumes finite (carrier R)
    shows finite-field R
  <proof>

  lemma (in domain) finite-domain-units:
    assumes finite (carrier R)
    shows Units R = carrier R - {0} (is ?lhs = ?rhs)
  <proof>

```

The following theorem can be found in Lidl and Niederreiter [4, Theorem 1.31].

theorem *finite-domains-are-fields*:
assumes *domain R*
assumes *finite (carrier R)*
shows *finite-field R*
 \langle *proof* \rangle

definition *zfact-iso* :: *nat* \Rightarrow *nat* \Rightarrow *int set* **where**
zfact-iso *p k* = *Idl_Z {int p} +>_Z (int k)*

context
fixes *n* :: *nat*
assumes *n-gt-0*: *n > 0*
begin

private abbreviation *I* **where** *I* \equiv *Idl_Z {int n}*

private lemma *ideal-I*: *ideal I Z*
 \langle *proof* \rangle

lemma *int-cosetI*:
assumes *u mod (int n) = v mod (int n)*
shows *Idl_Z {int n} +>_Z u = Idl_Z {int n} +>_Z v*

 \langle *proof* \rangle

lemma *zfact-iso-inj*:
*inj-on (zfact-iso n) {..*n*}*
 \langle *proof* \rangle

lemma *zfact-iso-ran*:
*zfact-iso n ' {..*n*} = carrier (ZFact (int n))*
 \langle *proof* \rangle

lemma *zfact-iso-bij*:
*bij-betw (zfact-iso n) {..*n*} (carrier (ZFact (int n)))*
 \langle *proof* \rangle

lemma *card-zfact-carr*: *card (carrier (ZFact (int n))) = n*
 \langle *proof* \rangle

lemma *fin-zfact*: *finite (carrier (ZFact (int n)))*
 \langle *proof* \rangle

end

lemma *zfact-prime-is-finite-field*:
assumes *Factorial-Ring.prime p*
shows *finite-field (ZFact (int p))*
 \langle *proof* \rangle

definition *int-embed* :: $- \Rightarrow \text{int} \Rightarrow -$ **where**
int-embed R $k = \text{add-pow } R$ k $\mathbf{1}_R$

lemma (**in** *ring*) *add-pow-consistent*:
fixes $i :: \text{int}$
assumes *subring* K R
assumes $k \in K$
shows $\text{add-pow } R$ i $k = \text{add-pow } (R$ (\mid *carrier* := K \mid)) i k
(is $?lhs = ?rhs$)
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-consistent*:
assumes *subring* K R
shows $\text{int-embed } R$ $i = \text{int-embed } (R$ (\mid *carrier* := K \mid)) i
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-closed*:
int-embed R $k \in \text{carrier } R$
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-range*:
assumes *subring* K R
shows $\text{int-embed } R$ $k \in K$
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-zero*:
int-embed R $0 = \mathbf{0}_R$
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-one*:
int-embed R $1 = \mathbf{1}_R$
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-add*:
int-embed R $(x+y) = \text{int-embed } R$ $x \oplus_R \text{int-embed } R$ y
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-inv*:
int-embed R $(-x) = \ominus_R \text{int-embed } R$ x **(is** $?lhs = ?rhs$)
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-diff*:
int-embed R $(x-y) = \text{int-embed } R$ $x \ominus_R \text{int-embed } R$ y
(is $?lhs = ?rhs$)
 $\langle \text{proof} \rangle$

lemma (**in** *ring*) *int-embed-mult-aux*:
int-embed R $(x*\text{int } y) = \text{int-embed } R$ $x \otimes \text{int-embed } R$ y
 $\langle \text{proof} \rangle$

lemma (in ring) *int-embed-mult*:
*int-embed R (x*y) = int-embed R x* \otimes_R *int-embed R y*
 <proof>

lemma (in ring) *int-embed-ring-hom*:
ring-hom-ring int-ring R (int-embed R)
 <proof>

abbreviation *char-subring where*
char-subring R \equiv *int-embed R* ‘UNIV

definition *char where*
char R = card (char-subring R)

This is a non-standard definition for the characteristic of a ring. Commonly [4, Definition 1.43] it is defined to be the smallest natural number n such that n -times repeated addition of any number is zero. If no such number exists then it is defined to be 0. In the case of rings with unit elements — not that the locale *Ring.ring* requires unit elements — the above definition can be simplified to the number of times the unit elements needs to be repeatedly added to reach 0.

The following three lemmas imply that the definition of the characteristic here coincides with the latter definition.

lemma (in ring) *char-bound*:
assumes $x > 0$
assumes *int-embed R (int x) = 0*
shows $\text{char } R \leq x$ $\text{char } R > 0$
 <proof>

lemma (in ring) *embed-char-eq-0*:
int-embed R (int (char R)) = 0
 <proof>

lemma (in ring) *embed-char-eq-0-iff*:
fixes $n :: \text{int}$
shows $\text{int-embed } R \ n = 0 \longleftrightarrow \text{char } R \ \text{dvd } n$
 <proof>

This result can be found in [4, Theorem 1.44].

lemma (in domain) *characteristic-is-prime*:
assumes $\text{char } R > 0$
shows *prime (char R)*
 <proof>

lemma (in ring) *char-ring-is-subring*:
subring (char-subring R) R

<proof>

lemma (in *cring*) *char-ring-is-subcring*:
 subcring (*char-subring* R) R
 <proof>

lemma (in *domain*) *char-ring-is-subdomain*:
 subdomain (*char-subring* R) R
 <proof>

lemma *image-set-eqI*:
 assumes $\bigwedge x. x \in A \implies f x \in B$
 assumes $\bigwedge x. x \in B \implies g x \in A \wedge f (g x) = x$
 shows $f ' A = B$
 <proof>

This is the binomial expansion theorem for commutative rings.

lemma (in *cring*) *binomial-expansion*:
 fixes $n :: nat$
 assumes [*simp*]: $x \in carrier\ R\ y \in carrier\ R$
 shows $(x \oplus y) [\wedge] n =$
 $(\bigoplus k \in \{..n\}. int-embed\ R\ (n\ choose\ k) \otimes x [\wedge] k \otimes y [\wedge] (n-k))$
 <proof>

lemma *bin-prime-factor*:
 assumes *prime* p
 assumes $k > 0\ k < p$
 shows $p\ dvd\ (p\ choose\ k)$
 <proof>

theorem (in *domain*) *freshmans-dream*:
 assumes $char\ R > 0$
 assumes [*simp*]: $x \in carrier\ R\ y \in carrier\ R$
 shows $(x \oplus y) [\wedge] (char\ R) = x [\wedge] char\ R \oplus y [\wedge] char\ R$
 (is ?lhs = ?rhs)
 <proof>

The following theorem is sometimes called Freshman's dream for obvious reasons, it can be found in Lidl and Niederreiter [4, Theorem 1.46].

lemma (in *domain*) *freshmans-dream-ext*:
 fixes m
 assumes $char\ R > 0$
 assumes [*simp*]: $x \in carrier\ R\ y \in carrier\ R$
 defines $n \equiv char\ R^{\wedge} m$
 shows $(x \oplus y) [\wedge] n = x [\wedge] n \oplus y [\wedge] n$
 (is ?lhs = ?rhs)
 <proof>

The following is a generalized version of the Frobenius homomorphism. The classic version of the theorem is the case where $k = 1$.

theorem (in domain) *frobenius-hom*:

assumes $\text{char } R > 0$

assumes $m = \text{char } R \wedge k$

shows $\text{ring-hom-cring } R R (\lambda x. x [\wedge] m)$

<proof>

lemma (in domain) *char-ring-is-subfield*:

assumes $\text{char } R > 0$

shows $\text{subfield } (\text{char-subring } R) R$

<proof>

lemma *card-lists-length-eq'*:

fixes $A :: 'a \text{ set}$

shows $\text{card } \{xs. \text{set } xs \subseteq A \wedge \text{length } xs = n\} = \text{card } A \wedge n$

<proof>

lemma (in ring) *card-span*:

assumes $\text{subfield } K R$

assumes $\text{independent } K w$

assumes $\text{set } w \subseteq \text{carrier } R$

shows $\text{card } (\text{Span } K w) = \text{card } K \wedge (\text{length } w)$

<proof>

lemma (in ring) *finite-carr-imp-char-ge-0*:

assumes $\text{finite } (\text{carrier } R)$

shows $\text{char } R > 0$

<proof>

lemma (in ring) *char-consistent*:

assumes $\text{subring } H R$

shows $\text{char } (R \upharpoonright \text{carrier } := H) = \text{char } R$

<proof>

lemma (in ring-hom-ring) *char-consistent*:

assumes $\text{inj-on } h (\text{carrier } R)$

shows $\text{char } R = \text{char } S$

<proof>

definition *char-iso* :: $- \Rightarrow \text{int set} \Rightarrow 'a$

where $\text{char-iso } R x = \text{the-elem } (\text{int-embed } R \text{ ' } x)$

The function $\text{char-iso } R$ denotes the isomorphism between $Z\text{Fact } (\text{int } (\text{char } R))$ and the characteristic subring.

lemma (in ring) *char-iso*: $\text{char-iso } R \in$

$\text{ring-iso } (Z\text{Fact } (\text{char } R)) (R \upharpoonright \text{carrier } := \text{char-subring } R)$

<proof>

The size of a finite field must be a prime power. This can be found in Ireland and Rosen [3, Proposition 7.1.3].

theorem (in *finite-field*) *finite-field-order*:

$\exists n. \text{order } R = \text{char } R \wedge n \wedge n > 0$

<proof>

end

4 Formal Derivatives

theory *Formal-Polynomial-Derivatives*

imports *HOL-Algebra.Polynomial-Divisibility Ring-Characteristic*

begin

definition *pderiv* (*pderiv*) **where**

$\text{pderiv}_R x = \text{ring.normalize } R ($

$\text{map } (\lambda i. \text{int-embed } R i \otimes_R \text{ring.coeff } R x i) (\text{rev } [1..<\text{length } x]))$

context *domain*

begin

lemma *coeff-range*:

assumes *subring* $K R$

assumes $f \in \text{carrier } (K[X])$

shows $\text{coeff } f i \in K$

<proof>

lemma *pderiv-carr*:

assumes *subring* $K R$

assumes $f \in \text{carrier } (K[X])$

shows $\text{pderiv } f \in \text{carrier } (K[X])$

<proof>

lemma *pderiv-coeff*:

assumes *subring* $K R$

assumes $f \in \text{carrier } (K[X])$

shows $\text{coeff } (\text{pderiv } f) k = \text{int-embed } R (\text{Suc } k) \otimes \text{coeff } f (\text{Suc } k)$

(**is** *?lhs = ?rhs*)

<proof>

lemma *pderiv-const*:

assumes *degree* $x = 0$

shows $\text{pderiv } x = \mathbf{0}_{K[X]}$

<proof>

lemma *pderiv-var*:

shows $\text{pderiv } X = \mathbf{1}_{K[X]}$

<proof>

lemma *pderiv-zero*:

shows $pderiv \mathbf{0}_{K[X]} = \mathbf{0}_{K[X]}$
<proof>

lemma *pderiv-add*:

assumes *subring* $K R$
assumes [*simp*]: $f \in carrier (K[X]) \ g \in carrier (K[X])$
shows $pderiv (f \oplus_{K[X]} g) = pderiv f \oplus_{K[X]} pderiv g$
(**is** ?lhs = ?rhs)
<proof>

lemma *pderiv-inv*:

assumes *subring* $K R$
assumes [*simp*]: $f \in carrier (K[X])$
shows $pderiv (\ominus_{K[X]} f) = \ominus_{K[X]} pderiv f$ (**is** ?lhs = ?rhs)
<proof>

lemma *coeff-mult*:

assumes *subring* $K R$
assumes $f \in carrier (K[X]) \ g \in carrier (K[X])$
shows $coeff (f \otimes_{K[X]} g) i =$
 $(\bigoplus k \in \{..i\}. (coeff f) k \otimes (coeff g) (i - k))$
<proof>

lemma *pderiv-mult*:

assumes *subring* $K R$
assumes [*simp*]: $f \in carrier (K[X]) \ g \in carrier (K[X])$
shows $pderiv (f \otimes_{K[X]} g) =$
 $pderiv f \otimes_{K[X]} g \oplus_{K[X]} f \otimes_{K[X]} pderiv g$
(**is** ?lhs = ?rhs)
<proof>

lemma *pderiv-pow*:

assumes $n > (0 :: nat)$
assumes *subring* $K R$
assumes [*simp*]: $f \in carrier (K[X])$
shows $pderiv (f [\wedge]_{K[X]} n) =$
 $int_embed (K[X]) n \otimes_{K[X]} f [\wedge]_{K[X]} (n-1) \otimes_{K[X]} pderiv f$
(**is** ?lhs = ?rhs)
<proof>

lemma *pderiv-var-pow*:

assumes $n > (0 :: nat)$
assumes *subring* $K R$
shows $pderiv (X [\wedge]_{K[X]} n) =$
 $int_embed (K[X]) n \otimes_{K[X]} X [\wedge]_{K[X]} (n-1)$

⟨proof⟩

lemma *int-embed-consistent-with-poly-of-const*:

assumes *subring* $K R$

shows *int-embed* $(K[X]) m = \text{poly-of-const} (\text{int-embed } R m)$

⟨proof⟩

end

end

5 Factorization into Monic Polynomials

theory *Monic-Polynomial-Factorization*

imports

Finite-Fields-Factorization-Ext

Formal-Polynomial-Derivatives

begin

hide-const *Factorial-Ring.multiplicity*

hide-const *Factorial-Ring.irreducible*

lemma (**in** *domain*) *finprod-mult-of*:

assumes *finite* A

assumes $\bigwedge x. x \in A \implies f x \in \text{carrier} (\text{mult-of } R)$

shows *finprod* $R f A = \text{finprod} (\text{mult-of } R) f A$

⟨proof⟩

lemma (**in** *ring*) *finite-poly*:

assumes *subring* $K R$

assumes *finite* K

shows

finite $\{f. f \in \text{carrier} (K[X]) \wedge \text{degree } f = n\}$ (**is finite** ? A)

finite $\{f. f \in \text{carrier} (K[X]) \wedge \text{degree } f \leq n\}$ (**is finite** ? B)

⟨proof⟩

definition *pmult* :: $- \Rightarrow 'a \text{ list} \Rightarrow 'a \text{ list} \Rightarrow \text{nat}$ (*pmulti*)

where *pmult* _{R} $d p = \text{multiplicity} (\text{mult-of} (\text{poly-ring } R)) d p$

definition *monic-poly* :: $- \Rightarrow 'a \text{ list} \Rightarrow \text{bool}$

where *monic-poly* $R f =$

$(f \neq [] \wedge \text{lead-coeff } f = \mathbf{1}_R \wedge f \in \text{carrier} (\text{poly-ring } R))$

definition *monic-irreducible-poly* **where**

monic-irreducible-poly $R f =$

$(\text{monic-poly } R f \wedge \text{pirreducible}_R (\text{carrier } R) f)$

abbreviation *m-i-p* $\equiv \text{monic-irreducible-poly}$

```

locale polynomial-ring = field +
  fixes K
  assumes polynomial-ring-assms: subfield K R
begin

```

```

lemma K-subring: subring K R
  ⟨proof⟩

```

```

abbreviation P where  $P \equiv K[X]$ 

```

This locale is used to specialize the following lemmas for a fixed coefficient ring. It can be introduced in a context as an interpretation to be able to use the following specialized lemmas. Because it is not (and should not) be introduced as a sublocale it has no lasting effect for the field locale itself.

lemmas

```

  poly-mult-lead-coeff = poly-mult-lead-coeff[OF K-subring]
and degree-add-distinct = degree-add-distinct[OF K-subring]
and coeff-add = coeff-add[OF K-subring]
and var-closed = var-closed[OF K-subring]
and degree-prod = degree-prod[OF K-subring]
and degree-pow = degree-pow[OF K-subring]
and pirreducible-degree = pirreducible-degree[OF polynomial-ring-assms]
and degree-one-imp-pirreducible =
  degree-one-imp-pirreducible[OF polynomial-ring-assms]
and var-pow-closed = var-pow-closed[OF K-subring]
and var-pow-carr = var-pow-carr[OF K-subring]
and univ-poly-a-inv-degree = univ-poly-a-inv-degree[OF K-subring]
and var-pow-degree = var-pow-degree[OF K-subring]
and pdivides-zero = pdivides-zero[OF K-subring]
and pdivides-imp-degree-le = pdivides-imp-degree-le[OF K-subring]
and var-carr = var-carr[OF K-subring]
and rupture-eq-0-iff = rupture-eq-0-iff[OF polynomial-ring-assms]
and rupture-is-field-iff-pirreducible =
  rupture-is-field-iff-pirreducible[OF polynomial-ring-assms]
and rupture-surj-hom = rupture-surj-hom[OF K-subring]
and canonical-embedding-ring-hom =
  canonical-embedding-ring-hom[OF K-subring]
and rupture-surj-norm-is-hom = rupture-surj-norm-is-hom[OF K-subring]
and rupture-surj-as-eval = rupture-surj-as-eval[OF K-subring]
and eval-cring-hom = eval-cring-hom[OF K-subring]
and coeff-range = coeff-range[OF K-subring]
and finite-poly = finite-poly[OF K-subring]
and int-embed-consistent-with-poly-of-const =
  int-embed-consistent-with-poly-of-const[OF K-subring]
and pderiv-var-pow = pderiv-var-pow[OF K-subring]
and pderiv-add = pderiv-add[OF K-subring]
and pderiv-inv = pderiv-inv[OF K-subring]
and pderiv-mult = pderiv-mult[OF K-subring]

```

and $pderiv-pow = pderiv-pow[OF - K-subring]$
and $pderiv-carr = pderiv-carr[OF K-subring]$

sublocale $p:principal-domain\ poly-ring\ R$
 $\langle proof \rangle$

end

context $field$
begin

interpretation $polynomial-ring\ R\ carrier\ R$
 $\langle proof \rangle$

lemma $pdivides-mult-r$:
assumes $a \in carrier\ (mult-of\ P)$
assumes $b \in carrier\ (mult-of\ P)$
assumes $c \in carrier\ (mult-of\ P)$
shows $a \otimes_P c\ pdivides\ b \otimes_P c \longleftrightarrow a\ pdivides\ b$
 $(is\ ?lhs \longleftrightarrow\ ?rhs)$
 $\langle proof \rangle$

lemma $lead-coeff-carr$:
assumes $x \in carrier\ (mult-of\ P)$
shows $lead-coeff\ x \in carrier\ R - \{0\}$
 $\langle proof \rangle$

lemma $lead-coeff-poly-of-const$:
assumes $r \neq 0$
shows $lead-coeff\ (poly-of-const\ r) = r$
 $\langle proof \rangle$

lemma $lead-coeff-mult$:
assumes $f \in carrier\ (mult-of\ P)$
assumes $g \in carrier\ (mult-of\ P)$
shows $lead-coeff\ (f \otimes_P g) = lead-coeff\ f \otimes lead-coeff\ g$
 $\langle proof \rangle$

lemma $monic-poly-carr$:
assumes $monic-poly\ R\ f$
shows $f \in carrier\ P$
 $\langle proof \rangle$

lemma $monic-poly-add-distinct$:
assumes $monic-poly\ R\ f$
assumes $g \in carrier\ P\ degree\ g < degree\ f$
shows $monic-poly\ R\ (f \oplus_P g)$
 $\langle proof \rangle$

lemma *monic-poly-one: monic-poly R 1_P*
⟨proof⟩

lemma *monic-poly-var: monic-poly R X*
⟨proof⟩

lemma *monic-poly-carr-2:*
 assumes *monic-poly R f*
 shows $f \in \text{carrier } (\text{mult-of } P)$
⟨proof⟩

lemma *monic-poly-mult:*
 assumes *monic-poly R f*
 assumes *monic-poly R g*
 shows *monic-poly R (f ⊗_P g)*
⟨proof⟩

lemma *monic-poly-pow:*
 assumes *monic-poly R f*
 shows *monic-poly R (f [^]_P (n::nat))*
⟨proof⟩

lemma *monic-poly-prod:*
 assumes *finite A*
 assumes $\bigwedge x. x \in A \implies \text{monic-poly } R (f x)$
 shows *monic-poly R (finprod P f A)*
⟨proof⟩

lemma *monic-poly-not-assoc:*
 assumes *monic-poly R f*
 assumes *monic-poly R g*
 assumes $f \sim_{(\text{mult-of } P)} g$
 shows $f = g$
⟨proof⟩

lemma *monic-poly-span:*
 assumes $x \in \text{carrier } (\text{mult-of } P) \text{ irreducible } (\text{mult-of } P) x$
 shows $\exists y. \text{monic-irreducible-poly } R y \wedge x \sim_{(\text{mult-of } P)} y$
⟨proof⟩

lemma *monic-polys-are-canonical-irreducibles:*
 canonical-irreducibles (mult-of P) {d. monic-irreducible-poly R d}
 (is canonical-irreducibles (mult-of P) ?S)
⟨proof⟩

lemma
 assumes *monic-poly R a*
 shows *factor-monic-poly:*
 $a = (\bigotimes_{p \in d} \text{monic-irreducible-poly } R p) \wedge \text{pmult } d a > 0$.

$d [\bigwedge]_P \text{pmult } d \ a) \text{ (is ?lhs = ?rhs)}$
and *factor-monic-poly-fin*:
 $\text{finite } \{d. \text{monic-irreducible-poly } R \ d \wedge \text{pmult } d \ a > 0\}$
 <proof>

lemma *degree-monic-poly'*:
assumes *monic-poly* $R \ f$
shows
 $\text{sum}' (\lambda d. \text{pmult } d \ f * \text{degree } d) \{d. \text{monic-irreducible-poly } R \ d\} =$
 $\text{degree } f$
 <proof>

lemma *monic-poly-min-degree*:
assumes *monic-irreducible-poly* $R \ f$
shows $\text{degree } f \geq 1$
 <proof>

lemma *degree-one-monic-poly*:
 $\text{monic-irreducible-poly } R \ f \wedge \text{degree } f = 1 \longleftrightarrow$
 $(\exists x \in \text{carrier } R. f = [\mathbf{1}, \ominus x])$
 <proof>

lemma *multiplicity-ge-iff*:
assumes *monic-irreducible-poly* $R \ d$
assumes $f \in \text{carrier } P - \{\mathbf{0}_P\}$
shows $\text{pmult } d \ f \geq k \longleftrightarrow d [\bigwedge]_P k \text{pdivides } f$
 <proof>

lemma *multiplicity-ge-1-iff-pdivides*:
assumes *monic-irreducible-poly* $R \ d \ f \in \text{carrier } P - \{\mathbf{0}_P\}$
shows $\text{pmult } d \ f \geq 1 \longleftrightarrow d \text{pdivides } f$
 <proof>

lemma *divides-monic-poly*:
assumes *monic-poly* $R \ f \ \text{monic-poly } R \ g$
assumes $\bigwedge d. \text{monic-irreducible-poly } R \ d$
 $\implies \text{pmult } d \ f \leq \text{pmult } d \ g$
shows $f \text{pdivides } g$
 <proof>

end

lemma *monic-poly-hom*:
assumes *monic-poly* $R \ f$
assumes $h \in \text{ring-iso } R \ S \ \text{domain } R \ \text{domain } S$
shows *monic-poly* $S \ (\text{map } h \ f)$
 <proof>

lemma *monic-irreducible-poly-hom*:


```

assumes monic-irreducible-poly  $R$   $f$ 
assumes  $h \in \text{ring-iso } R \ S$  domain  $R$  domain  $S$ 
shows monic-irreducible-poly  $S$  (map  $h$   $f$ )
⟨proof⟩

end

```

6 Counting Irreducible Polynomials

6.1 The polynomial $X^n - X$

theory *Card-Irreducible-Polynomials-Aux*

imports

HOL-Algebra.Multiplicative-Group

Formal-Polynomial-Derivatives

Monic-Polynomial-Factorization

begin

lemma (*in domain*)

assumes *subfield* K R

assumes $f \in \text{carrier } (K[X])$ *degree* $f > 0$

shows *embed-inj*: *inj-on* (*rupture-surj* K $f \circ \text{poly-of-const}$) K

and *rupture-order*: *order* (*Rupt* K f) = *card* $K^{\widehat{\text{degree } f}}$

and *rupture-char*: *char* (*Rupt* K f) = *char* R

⟨*proof*⟩

definition *gauss-poly where*

gauss-poly K $n = X_K [\wedge]_{\text{poly-ring } K} (n::\text{nat}) \ominus_{\text{poly-ring } K} X_K$

context *field*

begin

interpretation *polynomial-ring* R *carrier* R

⟨*proof*⟩

The following lemma can be found in Ireland and Rosen [3, §7.1, Lemma 2].

lemma *gauss-poly-div-gauss-poly-iff-1*:

fixes l $m :: \text{nat}$

assumes $l > 0$

shows $(X [\wedge]_P l \ominus_P \mathbf{1}_P)$ *pdivides* $(X [\wedge]_P m \ominus_P \mathbf{1}_P) \iff l \text{ dvd } m$

(*is ?lhs* \iff *?rhs*)

⟨*proof*⟩

lemma *gauss-poly-factor*:

assumes $n > 0$

shows *gauss-poly* R $n = (X [\wedge]_P (n-1) \ominus_P \mathbf{1}_P) \otimes_P X$ (*is* $- =$ *?rhs*)

⟨*proof*⟩

lemma *var-neq-zero*: $X \neq \mathbf{0}_P$
 ⟨proof⟩

lemma *var-pow-eq-one-iff*: $X [\wedge]_P k = \mathbf{1}_P \longleftrightarrow k = (0::nat)$
 ⟨proof⟩

lemma *gauss-poly-carr*: $gauss-poly\ R\ n \in carrier\ P$
 ⟨proof⟩

lemma *gauss-poly-degree*:
 assumes $n > 1$
 shows $degree\ (gauss-poly\ R\ n) = n$
 ⟨proof⟩

lemma *gauss-poly-not-zero*:
 assumes $n > 1$
 shows $gauss-poly\ R\ n \neq \mathbf{0}_P$
 ⟨proof⟩

lemma *gauss-poly-monic*:
 assumes $n > 1$
 shows $monic-poly\ R\ (gauss-poly\ R\ n)$
 ⟨proof⟩

lemma *geom-nat*:
 fixes $q :: nat$
 fixes $x :: - :: \{comm-ring, monoid-mult\}$
 shows $(x-1) * (\sum i \in \{..<q\}. x^i) = x^q - 1$
 ⟨proof⟩

The following lemma can be found in Ireland and Rosen [3, §7.1, Lemma 3].

lemma *gauss-poly-div-gauss-poly-iff-2*:
 fixes $a :: int$
 fixes $l\ m :: nat$
 assumes $l > 0\ a > 1$
 shows $(a^l - 1) dvd (a^m - 1) \longleftrightarrow l\ dvd\ m$
 (is ?lhs \longleftrightarrow ?rhs)
 ⟨proof⟩

lemma *gauss-poly-div-gauss-poly-iff*:
 assumes $m > 0\ n > 0\ a > 1$
 shows $gauss-poly\ R\ (a^n) pdivides_R\ gauss-poly\ R\ (a^m)$
 $\longleftrightarrow n\ dvd\ m$ (is ?lhs=?rhs)
 ⟨proof⟩

end

context *finite-field*

begin

interpretation *polynomial-ring R carrier R*
⟨*proof*⟩

lemma *div-gauss-poly-iff:*

assumes $n > 0$

assumes *monic-irreducible-poly R f*

shows $f \text{ pdivides}_R \text{ gauss-poly } R \text{ (order } R^{\wedge}n) \longleftrightarrow \text{degree } f \text{ dvd } n$

⟨*proof*⟩

lemma *gauss-poly-splitted:*

splitted (gauss-poly R (order R))

⟨*proof*⟩

The following lemma, for the case when R is a simple prime field, can be found in Ireland and Rosen [3, §7.1, Theorem 2]. Here the result is verified even for arbitrary finite fields.

lemma *multiplicity-of-factor-of-gauss-poly:*

assumes $n > 0$

assumes *monic-irreducible-poly R f*

shows

$\text{pmult}_R f \text{ (gauss-poly } R \text{ (order } R^{\wedge}n)) = \text{of-bool (degree } f \text{ dvd } n)$

⟨*proof*⟩

The following lemma, for the case when R is a simple prime field, can be found in Ireland and Rosen [3, §7.1, Corollary 1]. Here the result is verified even for arbitrary finite fields.

lemma *card-irred-aux:*

assumes $n > 0$

shows $\text{order } R^{\wedge}n = (\sum d \mid d \text{ dvd } n. d *$

$\text{card } \{f. \text{monic-irreducible-poly } R f \wedge \text{degree } f = d\})$

(**is** *?lhs = ?rhs*)

⟨*proof*⟩

end

end

6.2 Gauss Formula

theory *Card-Irreducible-Polynomials*

imports

Dirichlet-Series.Moebius-Mu

Card-Irreducible-Polynomials-Aux

begin

hide-const *Polynomial.order*

The following theorem is a slightly generalized form of the formula discovered by Gauss for the number of monic irreducible polynomials over a finite field. He originally verified the result for the case when R is a simple prime field. The version of the formula here for the case where R may be an arbitrary finite field can be found in Chebolu and Mináč [1].

theorem (in *finite-field*) *card-irred*:

assumes $n > 0$

shows $n * \text{card} \{f. \text{monic-irreducible-poly } R f \wedge \text{degree } f = n\} =$
 $(\sum d \mid d \text{ dvd } n. \text{moebius-mu } d * (\text{order } R^{(n \text{ div } d)}))$

(**is** ?lhs = ?rhs)

⟨proof⟩

In the following an explicit analytic lower bound for the cardinality of monic irreducible polynomials is shown, with which existence follows. This part deviates from the classic approach, where existence is verified using a divisibility argument. The reason for the deviation is that an analytic bound can also be used to estimate the runtime of a randomized algorithm selecting an irreducible polynomial, by randomly sampling monic polynomials.

lemma (in *finite-field*) *card-irred-1*:

$\text{card} \{f. \text{monic-irreducible-poly } R f \wedge \text{degree } f = 1\} = \text{order } R$

⟨proof⟩

lemma (in *finite-field*) *card-irred-2*:

$\text{real} (\text{card} \{f. \text{monic-irreducible-poly } R f \wedge \text{degree } f = 2\}) =$
 $(\text{real} (\text{order } R)^2 - \text{order } R) / 2$

⟨proof⟩

lemma (in *finite-field*) *card-irred-gt-2*:

assumes $n > 2$

shows $\text{real} (\text{order } R)^n / (2 * \text{real } n) \leq$

$\text{card} \{f. \text{monic-irreducible-poly } R f \wedge \text{degree } f = n\}$

(**is** ?lhs \leq ?rhs)

⟨proof⟩

lemma (in *finite-field*) *exist-irred*:

assumes $n > 0$

obtains f **where** *monic-irreducible-poly* $R f$ *degree* $f = n$

⟨proof⟩

theorem *existence*:

assumes $n > 0$

assumes *Factorial-Ring.prime* p

shows $\exists (F:: \text{int set list set ring}). \text{finite-field } F \wedge \text{order } F = p^n$

⟨proof⟩

end

7 Isomorphism between Finite Fields

theory *Finite-Fields-Isomorphic*

imports

Card-Irreducible-Polynomials

begin

lemma (in *finite-field*) *eval-on-root-is-iso*:

defines $p \equiv \text{char } R$

assumes $f \in \text{carrier } (\text{poly-ring } (\text{ZFact } p))$

assumes $\text{pirreducible}_{(\text{ZFact } p)} (\text{carrier } (\text{ZFact } p)) f$

assumes $\text{order } R = p^{\wedge} \text{degree } f$

assumes $x \in \text{carrier } R$

assumes $\text{eval } (\text{map } (\text{char-iso } R) f) x = \mathbf{0}$

shows $\text{ring-hom-ring } (\text{Rupt}_{(\text{ZFact } p)} (\text{carrier } (\text{ZFact } p)) f) R$

$(\lambda g. \text{the-elem } ((\lambda g'. \text{eval } (\text{map } (\text{char-iso } R) g') x) 'g))$

$\langle \text{proof} \rangle$

lemma (in *domain*) *pdivides-consistent*:

assumes $\text{subfield } K R f \in \text{carrier } (K[X]) g \in \text{carrier } (K[X])$

shows $f \text{ pdivides } g \longleftrightarrow f \text{ pdivides } R \ (\text{carrier} := K) g$

$\langle \text{proof} \rangle$

lemma (in *finite-field*) *find-root*:

assumes $\text{subfield } K R$

assumes $\text{monic-irreducible-poly } (R \ (\text{carrier} := K)) f$

assumes $\text{order } R = \text{card } K^{\wedge} \text{degree } f$

obtains x **where** $\text{eval } f x = \mathbf{0} x \in \text{carrier } R$

$\langle \text{proof} \rangle$

lemma (in *finite-field*) *find-iso-from-zfact*:

defines $p \equiv \text{int } (\text{char } R)$

assumes $\text{monic-irreducible-poly } (\text{ZFact } p) f$

assumes $\text{order } R = \text{char } R^{\wedge} \text{degree } f$

shows $\exists \varphi. \varphi \in \text{ring-iso } (\text{Rupt}_{(\text{ZFact } p)} (\text{carrier } (\text{ZFact } p)) f) R$

$\langle \text{proof} \rangle$

theorem *uniqueness*:

assumes $\text{finite-field } F_1$

assumes $\text{finite-field } F_2$

assumes $\text{order } F_1 = \text{order } F_2$

shows $F_1 \simeq F_2$

$\langle \text{proof} \rangle$

end

References

- [1] S. K. Chebolu and J. Mináč. Counting irreducible polynomials over finite fields using the inclusion-exclusion principle. *Mathematics Magazine*, 84:369 – 371, 2010.
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- [3] K. Ireland and M. Rosen. *A classical introduction to modern number theory*, volume 84 of *Graduate texts in mathematics*. Springer, 1982.
- [4] R. Lidl and H. Niederreiter. *Introduction to Finite Fields and Their Applications*. Cambridge University Press, USA, 1986.