

Euler's Partition Theorem

Lukas Bulwahn

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Abstract

Euler's Partition Theorem states that the number of partitions with only distinct parts is equal to the number of partitions with only odd parts. The combinatorial proof follows John Harrison's pre-existing HOL Light formalization [1]. To understand the rough idea of the proof, I read the lecture notes of the MIT course 18.312 on Algebraic Combinatorics [2] by Gregg Musiker. This theorem is the 45th theorem of the Top 100 Theorems list.

Contents

1 (Finite) multisets	3
1.1 The type of multisets	3
1.2 Representing multisets	3
1.3 Basic operations	5
1.3.1 Conversion to set and membership	5
1.3.2 Union	7
1.3.3 Difference	8
1.3.4 Equality of multisets	9
1.3.5 Pointwise ordering induced by count	11
1.3.6 Intersection and bounded union	14
1.3.7 Additional intersection facts	15
1.3.8 Additional bounded union facts	17
1.3.9 Subset is an order	18
1.4 Replicate and repeat operations	18
1.4.1 Simprocs	19
1.4.2 Conditionally complete lattice	20
1.4.3 Filter (with comprehension syntax)	22
1.4.4 Size	23
1.5 Induction and case splits	25
1.5.1 Strong induction and subset induction for multisets	25
1.6 The fold combinator	26
1.7 Image	27
1.8 Further conversions	29

1.9	More properties of the replicate and repeat operations	32
1.10	Big operators	33
1.11	Alternative representations	39
1.11.1	Lists	39
1.12	The multiset order	41
1.12.1	Well-foundedness	41
1.12.2	Closure-free presentation	42
1.13	The multiset extension is cancellative for multiset union . . .	42
1.14	Quasi-executable version of the multiset extension	43
1.14.1	Partial-order properties	43
1.14.2	Monotonicity of multiset union	44
1.14.3	Termination proofs with multiset orders	44
1.15	Legacy theorem bindings	45
1.16	Naive implementation using lists	46
1.17	BNF setup	49
1.18	Size setup	51
2	Additions to Isabelle’s Main Theories	51
2.1	Addition to Finite-Set Theory	52
2.2	Additions to Groups-Big Theory	52
2.3	Addition to Set-Interval Theory	52
2.4	Additions to Multiset Theory	52
3	Number Partitions	53
3.1	Number Partitions as $nat \Rightarrow nat$ Functions	53
3.2	Bounds and Finiteness of Number Partitions	53
3.3	Operations of Number Partitions	54
3.4	Number Partitions as Multisets on Natural Numbers	55
3.4.1	Relationship to Definition on Functions	55
4	Euler’s Partition Theorem	56
4.1	Preliminaries	56
4.1.1	Additions to Divides Theory	56
4.1.2	Additions to Groups-Big Theory	56
4.1.3	Additions to Set-Interval Theory	56
4.1.4	Additions to Nat Theory or Power Theory	56
4.1.5	Additions to Finite-Set Theory	56
4.2	Binary Encoding of Natural Numbers	57
4.3	Decomposition of a Number into a Power of Two and an Odd Number	58
4.4	Partitions With Only Distinct and Only Odd Parts	59
4.5	Euler’s Partition Theorem	60

1 (Finite) multisets

```
theory Multiset  
imports Main  
begin
```

1.1 The type of multisets

```
definition multiset = {f :: 'a ⇒ nat. finite {x. f x > 0}}
```

```
typedef 'a multiset = multiset :: ('a ⇒ nat) set  
morphisms count Abs-multiset  
  <proof>
```

```
setup-lifting type-definition-multiset
```

```
lemma multiset-eq-iff:  $M = N \longleftrightarrow (\forall a. \text{count } M \ a = \text{count } N \ a)$   
  <proof>
```

```
lemma multiset-eqI:  $(\bigwedge x. \text{count } A \ x = \text{count } B \ x) \implies A = B$   
  <proof>
```

Preservation of the representing set *multiset*.

```
lemma const0-in-multiset:  $(\lambda a. 0) \in \text{multiset}$   
  <proof>
```

```
lemma only1-in-multiset:  $(\lambda b. \text{if } b = a \text{ then } n \text{ else } 0) \in \text{multiset}$   
  <proof>
```

```
lemma union-preserves-multiset:  $M \in \text{multiset} \implies N \in \text{multiset} \implies (\lambda a. M \ a + N \ a) \in \text{multiset}$   
  <proof>
```

```
lemma diff-preserves-multiset:  
  assumes  $M \in \text{multiset}$   
  shows  $(\lambda a. M \ a - N \ a) \in \text{multiset}$   
  <proof>
```

```
lemma filter-preserves-multiset:  
  assumes  $M \in \text{multiset}$   
  shows  $(\lambda x. \text{if } P \ x \text{ then } M \ x \text{ else } 0) \in \text{multiset}$   
  <proof>
```

```
lemmas in-multiset = const0-in-multiset only1-in-multiset  
  union-preserves-multiset diff-preserves-multiset filter-preserves-multiset
```

1.2 Representing multisets

Multiset enumeration

instantiation *multiset* :: (type) cancel-comm-monoid-add
begin

lift-definition *zero-multiset* :: 'a multiset is $\lambda a. 0$
<proof>

abbreviation *Mempty* :: 'a multiset ({#}) **where**
Mempty $\equiv 0$

lift-definition *plus-multiset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow 'a multiset is $\lambda M N.$
($\lambda a. M a + N a$)
<proof>

lift-definition *minus-multiset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow 'a multiset is λM
 $N. \lambda a. M a - N a$
<proof>

instance
<proof>

end

context
begin

qualified definition *is-empty* :: 'a multiset \Rightarrow bool **where**
[code-abbrev]: *is-empty* $A \longleftrightarrow A = \{\#\}$

end

lemma *add-mset-in-multiset*:
assumes $M: \langle M \in \text{multiset} \rangle$
shows $\langle \lambda b. \text{if } b = a \text{ then } \text{Suc } (M b) \text{ else } M b \rangle \in \text{multiset} \rangle$
<proof>

lift-definition *add-mset* :: 'a \Rightarrow 'a multiset \Rightarrow 'a multiset is
 $\lambda a M b. \text{if } b = a \text{ then } \text{Suc } (M b) \text{ else } M b$
<proof>

syntax
-multiset :: args \Rightarrow 'a multiset ({#(-)#})

translations
 $\{\#x, xs\# \} == \text{CONST } \text{add-mset } x \ \{\#xs\# \}$
 $\{\#x\# \} == \text{CONST } \text{add-mset } x \ \{\# \}$

lemma *count-empty* [simp]: count {#} a = 0
<proof>

lemma *count-add-mset* [simp]:

count (*add-mset* *b* *A*) *a* = (if *b* = *a* then *Suc* (*count* *A* *a*) else *count* *A* *a*)
<proof>

lemma *count-single*: *count* {#*b*#} *a* = (if *b* = *a* then 1 else 0)
<proof>

lemma
add-mset-not-empty [*simp*]: $\langle \text{add-mset } a \ A \neq \{\#\} \rangle$ **and**
empty-not-add-mset [*simp*]: $\{\#\} \neq \text{add-mset } a \ A$
<proof>

lemma *add-mset-add-mset-same-iff* [*simp*]:
add-mset *a* *A* = *add-mset* *a* *B* \longleftrightarrow *A* = *B*
<proof>

lemma *add-mset-commute*:
add-mset *x* (*add-mset* *y* *M*) = *add-mset* *y* (*add-mset* *x* *M*)
<proof>

1.3 Basic operations

1.3.1 Conversion to set and membership

definition *set-mset* :: 'a multiset \Rightarrow 'a set
where *set-mset* *M* = {*x*. *count* *M* *x* > 0}

abbreviation *Melem* :: 'a \Rightarrow 'a multiset \Rightarrow bool
where *Melem* *a* *M* \equiv *a* \in *set-mset* *M*

notation
Melem (*op* \in #) **and**
Melem ((- / \in # -) [51, 51] 50)

notation (ASCII)
Melem (*op* :#) **and**
Melem ((- / :# -) [51, 51] 50)

abbreviation *not-Melem* :: 'a \Rightarrow 'a multiset \Rightarrow bool
where *not-Melem* *a* *M* \equiv *a* \notin *set-mset* *M*

notation
not-Melem (*op* \notin #) **and**
not-Melem ((- / \notin # -) [51, 51] 50)

notation (ASCII)
not-Melem (*op* \sim :#) **and**
not-Melem ((- / \sim :# -) [51, 51] 50)

context
begin

qualified abbreviation $Ball :: 'a\ multiset \Rightarrow ('a \Rightarrow bool) \Rightarrow bool$
where $Ball\ M \equiv Set.Ball\ (set-mset\ M)$

qualified abbreviation $Bex :: 'a\ multiset \Rightarrow ('a \Rightarrow bool) \Rightarrow bool$
where $Bex\ M \equiv Set.Bex\ (set-mset\ M)$

end

syntax

$-MBall \quad ::\ pttrn \Rightarrow 'a\ set \Rightarrow bool \Rightarrow bool \quad ((\exists\forall\ -\in\#\ -\ ./\ -)\ [0, 0, 10]\ 10)$
 $-MBex \quad ::\ pttrn \Rightarrow 'a\ set \Rightarrow bool \Rightarrow bool \quad ((\exists\exists\ -\in\#\ -\ ./\ -)\ [0, 0, 10]\ 10)$

syntax (ASCII)

$-MBall \quad ::\ pttrn \Rightarrow 'a\ set \Rightarrow bool \Rightarrow bool \quad ((\exists\forall\ -:\#\ -\ ./\ -)\ [0, 0, 10]\ 10)$
 $-MBex \quad ::\ pttrn \Rightarrow 'a\ set \Rightarrow bool \Rightarrow bool \quad ((\exists\exists\ -:\#\ -\ ./\ -)\ [0, 0, 10]\ 10)$

translations

$\forall x \in\# A. P \Rightarrow CONST\ Multiset.Ball\ A\ (\lambda x. P)$
 $\exists x \in\# A. P \Rightarrow CONST\ Multiset.Bex\ A\ (\lambda x. P)$

lemma *count-eq-zero-iff*:

$count\ M\ x = 0 \longleftrightarrow x \notin\# M$
 $\langle proof \rangle$

lemma *not-in-iff*:

$x \notin\# M \longleftrightarrow count\ M\ x = 0$
 $\langle proof \rangle$

lemma *count-greater-zero-iff* [simp]:

$count\ M\ x > 0 \longleftrightarrow x \in\# M$
 $\langle proof \rangle$

lemma *count-inI*:

assumes $count\ M\ x = 0 \Longrightarrow False$
shows $x \in\# M$
 $\langle proof \rangle$

lemma *in-countE*:

assumes $x \in\# M$
obtains n **where** $count\ M\ x = Suc\ n$
 $\langle proof \rangle$

lemma *count-greater-eq-Suc-zero-iff* [simp]:

$count\ M\ x \geq Suc\ 0 \longleftrightarrow x \in\# M$
 $\langle proof \rangle$

lemma *count-greater-eq-one-iff* [simp]:

$count\ M\ x \geq 1 \longleftrightarrow x \in\# M$

$\langle proof \rangle$

lemma *set-mset-empty* [simp]:

$set-mset \{\#\} = \{\}$

$\langle proof \rangle$

lemma *set-mset-single*:

$set-mset \{\#b\# \} = \{b\}$

$\langle proof \rangle$

lemma *set-mset-eq-empty-iff* [simp]:

$set-mset M = \{\} \longleftrightarrow M = \{\#\}$

$\langle proof \rangle$

lemma *finite-set-mset* [iff]:

$finite (set-mset M)$

$\langle proof \rangle$

lemma *set-mset-add-mset-insert* [simp]: $\langle set-mset (add-mset a A) = insert a (set-mset A) \rangle$

$\langle proof \rangle$

lemma *multiset-nonemptyE* [elim]:

assumes $A \neq \{\#\}$

obtains x **where** $x \in \# A$

$\langle proof \rangle$

1.3.2 Union

lemma *count-union* [simp]:

$count (M + N) a = count M a + count N a$

$\langle proof \rangle$

lemma *set-mset-union* [simp]:

$set-mset (M + N) = set-mset M \cup set-mset N$

$\langle proof \rangle$

lemma *union-mset-add-mset-left* [simp]:

$add-mset a A + B = add-mset a (A + B)$

$\langle proof \rangle$

lemma *union-mset-add-mset-right* [simp]:

$A + add-mset a B = add-mset a (A + B)$

$\langle proof \rangle$

lemma *add-mset-add-single*: $\langle add-mset a A = A + \{\#a\#\} \rangle$

$\langle proof \rangle$

1.3.3 Difference

instance *multiset* :: (type) comm-monoid-diff
⟨proof⟩

lemma *count-diff* [simp]:
 $\text{count } (M - N) a = \text{count } M a - \text{count } N a$
⟨proof⟩

lemma *add-mset-diff-bothsides*:
⟨add-mset a M - add-mset a A = M - A⟩
⟨proof⟩

lemma *in-diff-count*:
 $a \in\# M - N \longleftrightarrow \text{count } N a < \text{count } M a$
⟨proof⟩

lemma *count-in-diffI*:
assumes $\bigwedge n. \text{count } N x = n + \text{count } M x \implies \text{False}$
shows $x \in\# M - N$
⟨proof⟩

lemma *in-diff-countE*:
assumes $x \in\# M - N$
obtains n **where** $\text{count } M x = \text{Suc } n + \text{count } N x$
⟨proof⟩

lemma *in-diffD*:
assumes $a \in\# M - N$
shows $a \in\# M$
⟨proof⟩

lemma *set-mset-diff*:
 $\text{set-mset } (M - N) = \{a. \text{count } N a < \text{count } M a\}$
⟨proof⟩

lemma *diff-empty* [simp]: $M - \{\#\} = M \wedge \{\#\} - M = \{\#\}$
⟨proof⟩

lemma *diff-cancel*: $A - A = \{\#\}$
⟨proof⟩

lemma *diff-union-cancelR*: $M + N - N = (M :: 'a \text{ multiset})$
⟨proof⟩

lemma *diff-union-cancelL*: $N + M - N = (M :: 'a \text{ multiset})$
⟨proof⟩

lemma *diff-right-commute*:
fixes $M N Q :: 'a \text{ multiset}$

shows $M - N - Q = M - Q - N$
 ⟨proof⟩

lemma *diff-add*:
fixes $M N Q :: 'a \text{ multiset}$
shows $M - (N + Q) = M - N - Q$
 ⟨proof⟩

lemma *insert-DiffM* [simp]: $x \in\# M \implies \text{add-mset } x (M - \{x\}) = M$
 ⟨proof⟩

lemma *insert-DiffM2*: $x \in\# M \implies (M - \{x\}) + \{x\} = M$
 ⟨proof⟩

lemma *diff-union-swap*: $a \neq b \implies \text{add-mset } b (M - \{a\}) = \text{add-mset } b M - \{a\}$
 ⟨proof⟩

lemma *diff-add-mset-swap* [simp]: $b \notin\# A \implies \text{add-mset } b M - A = \text{add-mset } b (M - A)$
 ⟨proof⟩

lemma *diff-union-swap2* [simp]: $y \in\# M \implies \text{add-mset } x M - \{y\} = \text{add-mset } x (M - \{y\})$
 ⟨proof⟩

lemma *diff-diff-add-mset* [simp]: $(M :: 'a \text{ multiset}) - N - P = M - (N + P)$
 ⟨proof⟩

lemma *diff-union-single-conv*:
 $a \in\# J \implies I + J - \{a\} = I + (J - \{a\})$
 ⟨proof⟩

lemma *mset-add* [elim?]:
assumes $a \in\# A$
obtains B **where** $A = \text{add-mset } a B$
 ⟨proof⟩

lemma *union-iff*:
 $a \in\# A + B \iff a \in\# A \vee a \in\# B$
 ⟨proof⟩

1.3.4 Equality of multisets

lemma *single-eq-single* [simp]: $\{a\} = \{b\} \iff a = b$
 ⟨proof⟩

lemma *union-eq-empty* [iff]: $M + N = \{\#\} \iff M = \{\#\} \wedge N = \{\#\}$
 ⟨proof⟩

lemma *empty-eq-union* [iff]: $\{\#\} = M + N \longleftrightarrow M = \{\#\} \wedge N = \{\#\}$
 ⟨proof⟩

lemma *multi-self-add-other-not-self* [simp]: $M = \text{add-mset } x \ M \longleftrightarrow \text{False}$
 ⟨proof⟩

lemma *add-mset-remove-trivial* [simp]: $\langle \text{add-mset } x \ M - \{\#x\# \} = M \rangle$
 ⟨proof⟩

lemma *diff-single-trivial*: $\neg x \in\# \ M \implies M - \{\#x\# \} = M$
 ⟨proof⟩

lemma *diff-single-eq-union*: $x \in\# \ M \implies M - \{\#x\# \} = N \longleftrightarrow M = \text{add-mset } x \ N$
 ⟨proof⟩

lemma *union-single-eq-diff*: $\text{add-mset } x \ M = N \implies M = N - \{\#x\# \}$
 ⟨proof⟩

lemma *union-single-eq-member*: $\text{add-mset } x \ M = N \implies x \in\# \ N$
 ⟨proof⟩

lemma *add-mset-remove-trivial-If*:
 $\text{add-mset } a \ (N - \{\#a\# \}) = (\text{if } a \in\# \ N \text{ then } N \text{ else } \text{add-mset } a \ N)$
 ⟨proof⟩

lemma *add-mset-remove-trivial-eq*: $\langle N = \text{add-mset } a \ (N - \{\#a\# \}) \longleftrightarrow a \in\# \ N \rangle$
 ⟨proof⟩

lemma *union-is-single*:
 $M + N = \{\#a\# \} \longleftrightarrow M = \{\#a\# \} \wedge N = \{\#\} \vee M = \{\#\} \wedge N = \{\#a\# \}$
 (is ?lhs = ?rhs)
 ⟨proof⟩

lemma *single-is-union*: $\{\#a\# \} = M + N \longleftrightarrow \{\#a\# \} = M \wedge N = \{\#\} \vee M = \{\#\} \wedge \{\#a\# \} = N$
 ⟨proof⟩

lemma *add-eq-conv-diff*:
 $\text{add-mset } a \ M = \text{add-mset } b \ N \longleftrightarrow M = N \wedge a = b \vee M = \text{add-mset } b \ (N - \{\#a\# \}) \wedge N = \text{add-mset } a \ (M - \{\#b\# \})$
 (is ?lhs \longleftrightarrow ?rhs)

⟨proof⟩

lemma *add-mset-eq-single* [iff]: $\text{add-mset } b \ M = \{\#a\# \} \longleftrightarrow b = a \wedge M = \{\#\}$
 ⟨proof⟩

lemma *single-eq-add-mset* [iff]: $\{\#a\# \} = \text{add-mset } b \ M \longleftrightarrow b = a \wedge M = \{\#\}$
 ⟨proof⟩

lemma *insert-noteq-member*:
assumes BC : $\text{add-mset } b \ B = \text{add-mset } c \ C$
and $bnotc$: $b \neq c$
shows $c \in\# \ B$
 ⟨proof⟩

lemma *add-eq-conv-ex*:
 $(\text{add-mset } a \ M = \text{add-mset } b \ N) =$
 $(M = N \wedge a = b \vee (\exists K. M = \text{add-mset } b \ K \wedge N = \text{add-mset } a \ K))$
 ⟨proof⟩

lemma *multi-member-split*: $x \in\# \ M \implies \exists A. M = \text{add-mset } x \ A$
 ⟨proof⟩

lemma *multiset-add-sub-el-shuffle*:
assumes $c \in\# \ B$
and $b \neq c$
shows $\text{add-mset } b \ (B - \{\#c\# \}) = \text{add-mset } b \ B - \{\#c\# \}$
 ⟨proof⟩

lemma *add-mset-eq-singleton-iff*[iff]:
 $\text{add-mset } x \ M = \{\#y\# \} \longleftrightarrow M = \{\#\} \wedge x = y$
 ⟨proof⟩

1.3.5 Pointwise ordering induced by count

definition *subseteq-mset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow bool (**infix** $\subseteq\#$ 50)
where $A \subseteq\# \ B = (\forall a. \text{count } A \ a \leq \text{count } B \ a)$

definition *subset-mset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow bool (**infix** $\subset\#$ 50)
where $A \subset\# \ B = (A \subseteq\# \ B \wedge A \neq B)$

abbreviation (*input*) *supseteq-mset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow bool (**infix** $\supseteq\#$ 50)
where $\text{supseteq-mset } A \ B \equiv B \subseteq\# \ A$

abbreviation (*input*) *supset-mset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow bool (**infix** $\supset\#$ 50)
where $\text{supset-mset } A \ B \equiv B \subset\# \ A$

notation (*input*)
 subseq-mset (**infix** $\leq\#$ 50) **and**
 supseq-mset (**infix** $\geq\#$ 50)

notation (*ASCII*)

subseteq-mset (**infix** $\leq\#$ 50) **and**
subset-mset (**infix** $<\#$ 50) **and**
supseteq-mset (**infix** $\geq\#$ 50) **and**
supset-mset (**infix** $>\#$ 50)

interpretation *subset-mset*: *ordered-ab-semigroup-add-imp-le* $op + op - op \subseteq\#$
 $op \subseteq\#$
 ⟨*proof*⟩

interpretation *subset-mset*: *ordered-ab-semigroup-monoid-add-imp-le* $op + 0 op$
 $- op \leq\# op <\#$
 ⟨*proof*⟩

lemma *mset-subset-eqI*:
 $(\bigwedge a. \text{count } A a \leq \text{count } B a) \implies A \subseteq\# B$
 ⟨*proof*⟩

lemma *mset-subset-eq-count*:
 $A \subseteq\# B \implies \text{count } A a \leq \text{count } B a$
 ⟨*proof*⟩

lemma *mset-subset-eq-exists-conv*: $(A::'a \text{ multiset}) \subseteq\# B \longleftrightarrow (\exists C. B = A + C)$
 ⟨*proof*⟩

interpretation *subset-mset*: *ordered-cancel-comm-monoid-diff* $op + 0 op \leq\# op$
 $<\# op -$
 ⟨*proof*⟩

declare *subset-mset.add-diff-assoc*[*simp*] *subset-mset.add-diff-assoc2*[*simp*]

lemma *mset-subset-eq-mono-add-right-cancel*: $(A::'a \text{ multiset}) + C \subseteq\# B + C$
 $\longleftrightarrow A \subseteq\# B$
 ⟨*proof*⟩

lemma *mset-subset-eq-mono-add-left-cancel*: $C + (A::'a \text{ multiset}) \subseteq\# C + B \longleftrightarrow$
 $A \subseteq\# B$
 ⟨*proof*⟩

lemma *mset-subset-eq-mono-add*: $(A::'a \text{ multiset}) \subseteq\# B \implies C \subseteq\# D \implies A +$
 $C \subseteq\# B + D$
 ⟨*proof*⟩

lemma *mset-subset-eq-add-left*: $(A::'a \text{ multiset}) \subseteq\# A + B$
 ⟨*proof*⟩

lemma *mset-subset-eq-add-right*: $B \subseteq\# (A::'a \text{ multiset}) + B$
 ⟨*proof*⟩

lemma *single-subset-iff* [*simp*]:

$\{\#a\# \} \subseteq\# M \longleftrightarrow a \in\# M$
 $\langle proof \rangle$

lemma *mset-subset-eq-single*: $a \in\# B \implies \{\#a\# \} \subseteq\# B$
 $\langle proof \rangle$

lemma *mset-subset-eq-add-mset-cancel*: $\langle add-mset a A \subseteq\# add-mset a B \longleftrightarrow A \subseteq\# B \rangle$
 $\langle proof \rangle$

lemma *multiset-diff-union-assoc*:
fixes $A B C D :: 'a multiset$
shows $C \subseteq\# B \implies A + B - C = A + (B - C)$
 $\langle proof \rangle$

lemma *mset-subset-eq-multiset-union-diff-commute*:
fixes $A B C D :: 'a multiset$
shows $B \subseteq\# A \implies A - B + C = A + C - B$
 $\langle proof \rangle$

lemma *diff-subset-eq-self[simp]*:
 $(M :: 'a multiset) - N \subseteq\# M$
 $\langle proof \rangle$

lemma *mset-subset-eqD*:
assumes $A \subseteq\# B$ **and** $x \in\# A$
shows $x \in\# B$
 $\langle proof \rangle$

lemma *mset-subsetD*:
 $A \subset\# B \implies x \in\# A \implies x \in\# B$
 $\langle proof \rangle$

lemma *set-mset-mono*:
 $A \subseteq\# B \implies set-mset A \subseteq set-mset B$
 $\langle proof \rangle$

lemma *mset-subset-eq-insertD*:
 $add-mset x A \subseteq\# B \implies x \in\# B \wedge A \subset\# B$
 $\langle proof \rangle$

lemma *mset-subset-insertD*:
 $add-mset x A \subset\# B \implies x \in\# B \wedge A \subset\# B$
 $\langle proof \rangle$

lemma *mset-subset-of-empty[simp]*: $A \subset\# \{\#\} \longleftrightarrow False$
 $\langle proof \rangle$

lemma *empty-subset-add-mset[simp]*: $\{\#\} <\# add-mset x M$

<proof>

lemma *empty-le*: $\{\#\} \subseteq\# A$
<proof>

lemma *insert-subset-eq-iff*:
 $add_mset\ a\ A \subseteq\# B \longleftrightarrow a \in\# B \wedge A \subseteq\# B - \{\#a\}$
<proof>

lemma *insert-union-subset-iff*:
 $add_mset\ a\ A \subset\# B \longleftrightarrow a \in\# B \wedge A \subset\# B - \{\#a\}$
<proof>

lemma *subset-eq-diff-conv*:
 $A - C \subseteq\# B \longleftrightarrow A \subseteq\# B + C$
<proof>

lemma *multi-psub-of-add-self* [*simp*]: $A \subset\# add_mset\ x\ A$
<proof>

lemma *multi-psub-self*: $A \subset\# A = False$
<proof>

lemma *mset-subset-add-mset* [*simp*]: $add_mset\ x\ N \subset\# add_mset\ x\ M \longleftrightarrow N \subset\# M$
<proof>

lemma *mset-subset-diff-self*: $c \in\# B \implies B - \{\#c\} \subset\# B$
<proof>

lemma *Diff-eq-empty-iff-mset*: $A - B = \{\#\} \longleftrightarrow A \subseteq\# B$
<proof>

lemma *add-mset-subseteq-single-iff* [*iff*]: $add_mset\ a\ M \subseteq\# \{\#b\} \longleftrightarrow M = \{\#\} \wedge a = b$
<proof>

1.3.6 Intersection and bounded union

definition *inf-subset-mset* :: $'a\ multiset \Rightarrow 'a\ multiset \Rightarrow 'a\ multiset$ (**infixl** $\cap\#$ 70) **where**
multiset-inter-def: $inf_subset_mset\ A\ B = A - (A - B)$

interpretation *subset-mset*: *semilattice-inf* *inf-subset-mset* *op* $\subseteq\#$ *op* $\subset\#$
<proof>

definition *sup-subset-mset* :: $'a\ multiset \Rightarrow 'a\ multiset \Rightarrow 'a\ multiset$ (**infixl** $\cup\#$ 70)
where *sup-subset-mset* $A\ B = A + (B - A)$ — FIXME irregular fact name

interpretation *subset-mset*: *semilattice-sup sup-subset-mset* $op \subseteq\# op \subset\#$
 $\langle proof \rangle$

interpretation *subset-mset*: *bounded-lattice-bot* $op \cap\# op \subseteq\# op \subset\#$
 $op \cup\# \{\#\}$
 $\langle proof \rangle$

1.3.7 Additional intersection facts

lemma *mset-inter-count* [*simp*]:
fixes $A B :: 'a \text{ mset}$
shows $count (A \cap\# B) x = \min (count A x) (count B x)$
 $\langle proof \rangle$

lemma *set-mset-inter* [*simp*]:
 $set-mset (A \cap\# B) = set-mset A \cap set-mset B$
 $\langle proof \rangle$

lemma *diff-intersect-left-idem* [*simp*]:
 $M - M \cap\# N = M - N$
 $\langle proof \rangle$

lemma *diff-intersect-right-idem* [*simp*]:
 $M - N \cap\# M = M - N$
 $\langle proof \rangle$

lemma *mset-inter-single*[*simp*]: $a \neq b \implies \{\#a\} \cap\# \{\#b\} = \{\#\}$
 $\langle proof \rangle$

lemma *mset-union-diff-commute*:
assumes $B \cap\# C = \{\#\}$
shows $A + B - C = A - C + B$
 $\langle proof \rangle$

lemma *disjunct-not-in*:
 $A \cap\# B = \{\#\} \longleftrightarrow (\forall a. a \notin\# A \vee a \notin\# B)$ (**is** $?P \longleftrightarrow ?Q$)
 $\langle proof \rangle$

lemma *inter-mset-empty-distrib-right*: $A \cap\# (B + C) = \{\#\} \longleftrightarrow A \cap\# B = \{\#\} \wedge A \cap\# C = \{\#\}$
 $\langle proof \rangle$

lemma *inter-mset-empty-distrib-left*: $(A + B) \cap\# C = \{\#\} \longleftrightarrow A \cap\# C = \{\#\} \wedge B \cap\# C = \{\#\}$
 $\langle proof \rangle$

lemma *add-mset-inter-add-mset*[*simp*]:
 $add-mset a A \cap\# add-mset a B = add-mset a (A \cap\# B)$

<proof>

lemma *add-mset-disjoint* [*simp*]:

$add-mset\ a\ A\ \cap\# \ B = \{\#\} \longleftrightarrow a \notin\# \ B \wedge A \cap\# \ B = \{\#\}$
 $\{\#\} = add-mset\ a\ A\ \cap\# \ B \longleftrightarrow a \notin\# \ B \wedge \{\#\} = A \cap\# \ B$
<proof>

lemma *disjoint-add-mset* [*simp*]:

$B \cap\# \ add-mset\ a\ A = \{\#\} \longleftrightarrow a \notin\# \ B \wedge B \cap\# \ A = \{\#\}$
 $\{\#\} = A \cap\# \ add-mset\ b\ B \longleftrightarrow b \notin\# \ A \wedge \{\#\} = A \cap\# \ B$
<proof>

lemma *inter-add-left1*: $\neg x \in\# \ N \implies (add-mset\ x\ M) \cap\# \ N = M \cap\# \ N$

<proof>

lemma *inter-add-left2*: $x \in\# \ N \implies (add-mset\ x\ M) \cap\# \ N = add-mset\ x\ (M \cap\# \ (N - \{\#x\}))$

<proof>

lemma *inter-add-right1*: $\neg x \in\# \ N \implies N \cap\# \ (add-mset\ x\ M) = N \cap\# \ M$

<proof>

lemma *inter-add-right2*: $x \in\# \ N \implies N \cap\# \ (add-mset\ x\ M) = add-mset\ x\ ((N - \{\#x\}) \cap\# \ M)$

<proof>

lemma *disjunct-set-mset-diff*:

assumes $M \cap\# \ N = \{\#\}$
shows $set-mset\ (M - N) = set-mset\ M$
<proof>

lemma *at-most-one-mset-mset-diff*:

assumes $a \notin\# \ M - \{\#a\}$
shows $set-mset\ (M - \{\#a\}) = set-mset\ M - \{a\}$
<proof>

lemma *more-than-one-mset-mset-diff*:

assumes $a \in\# \ M - \{\#a\}$
shows $set-mset\ (M - \{\#a\}) = set-mset\ M$
<proof>

lemma *inter-iff*:

$a \in\# \ A \cap\# \ B \longleftrightarrow a \in\# \ A \wedge a \in\# \ B$
<proof>

lemma *inter-union-distrib-left*:

$A \cap\# \ B + C = (A + C) \cap\# \ (B + C)$
<proof>

lemma *inter-union-distrib-right*:

$$C + A \cap\# B = (C + A) \cap\# (C + B)$$

<proof>

lemma *inter-subset-eq-union*:

$$A \cap\# B \subseteq\# A + B$$

<proof>

1.3.8 Additional bounded union facts

lemma *sup-subset-mset-count* [*simp*]: — FIXME irregular fact name

$$\text{count } (A \cup\# B) x = \max (\text{count } A x) (\text{count } B x)$$

<proof>

lemma *set-mset-sup* [*simp*]:

$$\text{set-mset } (A \cup\# B) = \text{set-mset } A \cup \text{set-mset } B$$

<proof>

lemma *sup-union-left1* [*simp*]: $\neg x \in\# N \implies (\text{add-mset } x M) \cup\# N = \text{add-mset } x (M \cup\# N)$

<proof>

lemma *sup-union-left2*: $x \in\# N \implies (\text{add-mset } x M) \cup\# N = \text{add-mset } x (M \cup\# (N - \{\#x\}))$

<proof>

lemma *sup-union-right1* [*simp*]: $\neg x \in\# N \implies N \cup\# (\text{add-mset } x M) = \text{add-mset } x (N \cup\# M)$

<proof>

lemma *sup-union-right2*: $x \in\# N \implies N \cup\# (\text{add-mset } x M) = \text{add-mset } x ((N - \{\#x\}) \cup\# M)$

<proof>

lemma *sup-union-distrib-left*:

$$A \cup\# B + C = (A + C) \cup\# (B + C)$$

<proof>

lemma *union-sup-distrib-right*:

$$C + A \cup\# B = (C + A) \cup\# (C + B)$$

<proof>

lemma *union-diff-inter-eq-sup*:

$$A + B - A \cap\# B = A \cup\# B$$

<proof>

lemma *union-diff-sup-eq-inter*:

$$A + B - A \cup\# B = A \cap\# B$$

<proof>

lemma *add-mset-union*:

$\langle \text{add-mset } a \ A \cup\# \ \text{add-mset } a \ B = \text{add-mset } a \ (A \cup\# \ B) \rangle$
 $\langle \text{proof} \rangle$

1.3.9 Subset is an order

interpretation *subset-mset*: *order op* $\subseteq\#$ *op* $\subset\#$ $\langle \text{proof} \rangle$

1.4 Replicate and repeat operations

definition *replicate-mset* :: $\text{nat} \Rightarrow 'a \Rightarrow 'a \ \text{multiset}$ **where**
 $\text{replicate-mset } n \ x = (\text{add-mset } x \ \hat{\ } n) \ \{\#\}$

lemma *replicate-mset-0*[*simp*]: $\text{replicate-mset } 0 \ x = \{\#\}$
 $\langle \text{proof} \rangle$

lemma *replicate-mset-Suc* [*simp*]: $\text{replicate-mset } (\text{Suc } n) \ x = \text{add-mset } x \ (\text{replicate-mset } n \ x)$
 $\langle \text{proof} \rangle$

lemma *count-replicate-mset*[*simp*]: $\text{count } (\text{replicate-mset } n \ x) \ y = (\text{if } y = x \ \text{then } n \ \text{else } 0)$
 $\langle \text{proof} \rangle$

fun *repeat-mset* :: $\text{nat} \Rightarrow 'a \ \text{multiset} \Rightarrow 'a \ \text{multiset}$ **where**
 $\text{repeat-mset } 0 \ - = \{\#\} \mid$
 $\text{repeat-mset } (\text{Suc } n) \ A = A + \text{repeat-mset } n \ A$

lemma *count-repeat-mset* [*simp*]: $\text{count } (\text{repeat-mset } i \ A) \ a = i * \text{count } A \ a$
 $\langle \text{proof} \rangle$

lemma *repeat-mset-right* [*simp*]: $\text{repeat-mset } a \ (\text{repeat-mset } b \ A) = \text{repeat-mset } (a * b) \ A$
 $\langle \text{proof} \rangle$

lemma *left-diff-repeat-mset-distrib'*: $\langle \text{repeat-mset } (i - j) \ u = \text{repeat-mset } i \ u - \text{repeat-mset } j \ u \rangle$
 $\langle \text{proof} \rangle$

lemma *left-add-mult-distrib-mset*:

$\text{repeat-mset } i \ u + (\text{repeat-mset } j \ u + k) = \text{repeat-mset } (i+j) \ u + k$
 $\langle \text{proof} \rangle$

lemma *repeat-mset-distrib*:

$\text{repeat-mset } (m + n) \ A = \text{repeat-mset } m \ A + \text{repeat-mset } n \ A$
 $\langle \text{proof} \rangle$

lemma *repeat-mset-distrib2*[*simp*]:

$\text{repeat-mset } n \ (A + B) = \text{repeat-mset } n \ A + \text{repeat-mset } n \ B$

$\langle proof \rangle$

lemma *repeat-mset-replicate-mset*[simp]:
 $repeat\text{-}mset\ n\ \{\#a\#\} = replicate\text{-}mset\ n\ a$
 $\langle proof \rangle$

lemma *repeat-mset-distrib-add-mset*[simp]:
 $repeat\text{-}mset\ n\ (add\text{-}mset\ a\ A) = replicate\text{-}mset\ n\ a + repeat\text{-}mset\ n\ A$
 $\langle proof \rangle$

lemma *repeat-mset-empty*[simp]: $repeat\text{-}mset\ n\ \{\#\} = \{\#\}$
 $\langle proof \rangle$

1.4.1 Simprocs

lemma *mset-diff-add-eq1*:
 $j \leq (i::nat) \implies ((repeat\text{-}mset\ i\ u + m) - (repeat\text{-}mset\ j\ u + n)) = ((repeat\text{-}mset\ (i-j)\ u + m) - n)$
 $\langle proof \rangle$

lemma *mset-diff-add-eq2*:
 $i \leq (j::nat) \implies ((repeat\text{-}mset\ i\ u + m) - (repeat\text{-}mset\ j\ u + n)) = (m - (repeat\text{-}mset\ (j-i)\ u + n))$
 $\langle proof \rangle$

lemma *mset-eq-add-iff1*:
 $j \leq (i::nat) \implies (repeat\text{-}mset\ i\ u + m = repeat\text{-}mset\ j\ u + n) = (repeat\text{-}mset\ (i-j)\ u + m = n)$
 $\langle proof \rangle$

lemma *mset-eq-add-iff2*:
 $i \leq (j::nat) \implies (repeat\text{-}mset\ i\ u + m = repeat\text{-}mset\ j\ u + n) = (m = repeat\text{-}mset\ (j-i)\ u + n)$
 $\langle proof \rangle$

lemma *mset-subseteq-add-iff1*:
 $j \leq (i::nat) \implies (repeat\text{-}mset\ i\ u + m \subseteq\# repeat\text{-}mset\ j\ u + n) = (repeat\text{-}mset\ (i-j)\ u + m \subseteq\# n)$
 $\langle proof \rangle$

lemma *mset-subseteq-add-iff2*:
 $i \leq (j::nat) \implies (repeat\text{-}mset\ i\ u + m \subseteq\# repeat\text{-}mset\ j\ u + n) = (m \subseteq\# repeat\text{-}mset\ (j-i)\ u + n)$
 $\langle proof \rangle$

lemma *mset-subset-add-iff1*:
 $j \leq (i::nat) \implies (repeat\text{-}mset\ i\ u + m \subset\# repeat\text{-}mset\ j\ u + n) = (repeat\text{-}mset\ (i-j)\ u + m \subset\# n)$
 $\langle proof \rangle$

lemma *mset-subset-add-iff2*:

$i \leq (j::nat) \implies (repeat\text{-}mset\ i\ u + m \subset\# repeat\text{-}mset\ j\ u + n) = (m \subset\# repeat\text{-}mset\ (j-i)\ u + n)$
<proof>

<ML>

1.4.2 Conditionally complete lattice

instantiation *multiset* :: (type) *Inf*
begin

lift-definition *Inf-multiset* :: 'a multiset set \Rightarrow 'a multiset **is**

$\lambda A\ i.$ if $A = \{\}$ then 0 else *Inf* (($\lambda f.$ f i) 'A)
<proof>

instance *<proof>*

end

lemma *Inf-multiset-empty*: *Inf* $\{\}$ = $\{\#\}$
<proof>

lemma *count-Inf-multiset-nonempty*: $A \neq \{\} \implies count\ (Inf\ A)\ x = Inf\ ((\lambda X. count\ X\ x)\ 'A)$
<proof>

instantiation *multiset* :: (type) *Sup*
begin

definition *Sup-multiset* :: 'a multiset set \Rightarrow 'a multiset **where**

Sup-multiset A = (if $A \neq \{\}$ \wedge *subset-mset.bdd-above* A then
Abs-multiset ($\lambda i.$ *Sup* (($\lambda X.$ *count* X i) 'A)) else $\{\#\}$)

lemma *Sup-multiset-empty*: *Sup* $\{\}$ = $\{\#\}$
<proof>

lemma *Sup-multiset-unbounded*: $\neg subset\text{-}mset.\text{bdd-above}\ A \implies Sup\ A = \{\#\}$
<proof>

instance *<proof>*

end

lemma *bdd-above-multiset-imp-bdd-above-count*:

assumes *subset-mset.bdd-above* (A :: 'a multiset set)

shows *bdd-above* $((\lambda X. \text{count } X \ x) \ 'A)$
<proof>

lemma *bdd-above-multiset-imp-finite-support*:
assumes $A \neq \{\}$ *subset-mset.bdd-above* $(A :: 'a \text{ multiset set})$
shows *finite* $(\bigcup X \in A. \{x. \text{count } X \ x > 0\})$
<proof>

lemma *Sup-multiset-in-multiset*:
assumes $A \neq \{\}$ *subset-mset.bdd-above* A
shows $(\lambda i. \text{SUP } X:A. \text{count } X \ i) \in \text{multiset}$
<proof>

lemma *count-Sup-multiset-nonempty*:
assumes $A \neq \{\}$ *subset-mset.bdd-above* A
shows $\text{count } (\text{Sup } A) \ x = (\text{SUP } X:A. \text{count } X \ x)$
<proof>

interpretation *subset-mset: conditionally-complete-lattice* *Inf Sup op* $\cap\#$ *op* $\subseteq\#$
op $\subset\#$ *op* $\cup\#$
<proof>

lemma *set-mset-Inf*:
assumes $A \neq \{\}$
shows $\text{set-mset } (\text{Inf } A) = (\bigcap X \in A. \text{set-mset } X)$
<proof>

lemma *in-Inf-multiset-iff*:
assumes $A \neq \{\}$
shows $x \in\# \text{Inf } A \longleftrightarrow (\forall X \in A. x \in\# X)$
<proof>

lemma *in-Inf-multisetD*: $x \in\# \text{Inf } A \implies X \in A \implies x \in\# X$
<proof>

lemma *set-mset-Sup*:
assumes *subset-mset.bdd-above* A
shows $\text{set-mset } (\text{Sup } A) = (\bigcup X \in A. \text{set-mset } X)$
<proof>

lemma *in-Sup-multiset-iff*:
assumes *subset-mset.bdd-above* A
shows $x \in\# \text{Sup } A \longleftrightarrow (\exists X \in A. x \in\# X)$
<proof>

lemma *in-Sup-multisetD*:
assumes $x \in\# \text{Sup } A$
shows $\exists X \in A. x \in\# X$

<proof>

interpretation *subset-mset: distrib-lattice op* $\cap\#$ *op* $\subseteq\#$ *op* $\subset\#$ *op* $\cup\#$
<proof>

1.4.3 Filter (with comprehension syntax)

Multiset comprehension

lift-definition *filter-mset* :: ('a \Rightarrow bool) \Rightarrow 'a multiset \Rightarrow 'a multiset
is $\lambda P M. \lambda x. \text{if } P \ x \ \text{then } M \ x \ \text{else } 0$
<proof>

syntax (ASCII)

-MCollect :: ptrn \Rightarrow 'a multiset \Rightarrow bool \Rightarrow 'a multiset ((1{#- :# -./ -#}))

syntax

-MCollect :: ptrn \Rightarrow 'a multiset \Rightarrow bool \Rightarrow 'a multiset ((1{#- $\in\#$ -./ -#}))

translations

{#x $\in\#$ M. P#} == CONST *filter-mset* ($\lambda x. P$) M

lemma *count-filter-mset* [simp]:

count (*filter-mset* P M) a = (if P a then *count* M a else 0)

<proof>

lemma *set-mset-filter* [simp]:

set-mset (*filter-mset* P M) = {a \in *set-mset* M. P a}

<proof>

lemma *filter-empty-mset* [simp]: *filter-mset* P {#} = {#}

<proof>

lemma *filter-single-mset*: *filter-mset* P {#x#} = (if P x then {#x#} else {#})

<proof>

lemma *filter-union-mset* [simp]: *filter-mset* P (M + N) = *filter-mset* P M +
filter-mset P N

<proof>

lemma *filter-diff-mset* [simp]: *filter-mset* P (M - N) = *filter-mset* P M - *filter-mset*
P N

<proof>

lemma *filter-inter-mset* [simp]: *filter-mset* P (M $\cap\#$ N) = *filter-mset* P M $\cap\#$
filter-mset P N

<proof>

lemma *filter-sup-mset*[simp]: *filter-mset* P (A $\cup\#$ B) = *filter-mset* P A $\cup\#$ *filter-mset*
P B

<proof>

lemma *filter-mset-add-mset* [*simp*]:
 $filter\text{-}mset\ P\ (add\text{-}mset\ x\ A) =$
(if $P\ x$ *then* $add\text{-}mset\ x\ (filter\text{-}mset\ P\ A)$ *else* $filter\text{-}mset\ P\ A$)
 ⟨*proof*⟩

lemma *multiset-filter-subset*[*simp*]: $filter\text{-}mset\ f\ M \subseteq\# M$
 ⟨*proof*⟩

lemma *multiset-filter-mono*:
assumes $A \subseteq\# B$
shows $filter\text{-}mset\ f\ A \subseteq\# filter\text{-}mset\ f\ B$
 ⟨*proof*⟩

lemma *filter-mset-eq-conv*:
 $filter\text{-}mset\ P\ M = N \longleftrightarrow N \subseteq\# M \wedge (\forall b \in\# N. P\ b) \wedge (\forall a \in\# M - N. \neg P\ a)$
(is $?P \longleftrightarrow ?Q$)
 ⟨*proof*⟩

lemma *filter-filter-mset*: $filter\text{-}mset\ P\ (filter\text{-}mset\ Q\ M) = \{\#x \in\# M. Q\ x \wedge P\ x\}$
 ⟨*proof*⟩

lemma
 $filter\text{-}mset\ True$ [*simp*]: $\{\#y \in\# M. True\# \} = M$ **and**
 $filter\text{-}mset\ False$ [*simp*]: $\{\#y \in\# M. False\# \} = \{\#\}$
 ⟨*proof*⟩

1.4.4 Size

definition *wcount* **where** $wcount\ f\ M = (\lambda x. count\ M\ x * Suc\ (f\ x))$

lemma *wcount-union*: $wcount\ f\ (M + N)\ a = wcount\ f\ M\ a + wcount\ f\ N\ a$
 ⟨*proof*⟩

lemma *wcount-add-mset*:
 $wcount\ f\ (add\text{-}mset\ x\ M)\ a = (if\ x = a\ then\ Suc\ (f\ a)\ else\ 0) + wcount\ f\ M\ a$
 ⟨*proof*⟩

definition *size-multiset* :: $('a \Rightarrow nat) \Rightarrow 'a\ multiset \Rightarrow nat$ **where**
 $size\text{-}multiset\ f\ M = sum\ (wcount\ f\ M)\ (set\text{-}mset\ M)$

lemmas $size\text{-}multiset\text{-}eq = size\text{-}multiset\text{-}def$ [*unfolded* *wcount-def*]

instantiation *multiset* :: $(type)\ size$
begin

definition *size-multiset* **where**
 $size\text{-}multiset\text{-}overloaded\text{-}def: size\text{-}multiset = Multiset.size\text{-}multiset\ (\lambda_. 0)$
instance ⟨*proof*⟩

end

lemmas *size-multiset-overloaded-eq* =
size-multiset-overloaded-def[*THEN fun-cong, unfolded size-multiset-eq, simplified*]

lemma *size-multiset-empty* [*simp*]: *size-multiset* f $\{\#\}$ = 0
(*proof*)

lemma *size-empty* [*simp*]: *size* $\{\#\}$ = 0
(*proof*)

lemma *size-multiset-single* : *size-multiset* f $\{\#b\#\}$ = *Suc* (f b)
(*proof*)

lemma *size-single*: *size* $\{\#b\#\}$ = 1
(*proof*)

lemma *sum-wcount-Int*:
finite $A \implies \text{sum } (wcount\ f\ N) (A \cap \text{set-mset } N) = \text{sum } (wcount\ f\ N) A$
(*proof*)

lemma *size-multiset-union* [*simp*]:
size-multiset f ($M + N :: 'a$ *multiset*) = *size-multiset* f M + *size-multiset* f N
(*proof*)

lemma *size-multiset-add-mset* [*simp*]:
size-multiset f (*add-mset* a M) = *Suc* (f a) + *size-multiset* f M
(*proof*)

lemma *size-add-mset* [*simp*]: *size* (*add-mset* a A) = *Suc* (*size* A)
(*proof*)

lemma *size-union* [*simp*]: *size* ($M + N :: 'a$ *multiset*) = *size* M + *size* N
(*proof*)

lemma *size-multiset-eq-0-iff-empty* [*iff*]:
size-multiset f M = 0 $\iff M = \{\#\}$
(*proof*)

lemma *size-eq-0-iff-empty* [*iff*]: (*size* M = 0) = ($M = \{\#\}$)
(*proof*)

lemma *nonempty-has-size*: ($S \neq \{\#\}$) = ($0 < \text{size } S$)
(*proof*)

lemma *size-eq-Suc-imp-elem*: *size* M = *Suc* $n \implies \exists a. a \in\# M$
(*proof*)

lemma *size-eq-Suc-imp-eq-union*:
assumes $size\ M = Suc\ n$
shows $\exists a\ N. M = add\text{-}mset\ a\ N$
 $\langle proof \rangle$

lemma *size-mset-mono*:
fixes $A\ B :: 'a\ multiset$
assumes $A \subseteq\# B$
shows $size\ A \leq size\ B$
 $\langle proof \rangle$

lemma *size-filter-mset-lesseq[simp]*: $size\ (filter\text{-}mset\ f\ M) \leq size\ M$
 $\langle proof \rangle$

lemma *size-Diff-submset*:
 $M \subseteq\# M' \implies size\ (M' - M) = size\ M' - size\ (M :: 'a\ multiset)$
 $\langle proof \rangle$

1.5 Induction and case splits

theorem *multiset-induct* [*case-names empty add, induct type: multiset*]:
assumes *empty*: $P\ \{\#\}$
assumes *add*: $\bigwedge x\ M. P\ M \implies P\ (add\text{-}mset\ x\ M)$
shows $P\ M$
 $\langle proof \rangle$

lemma *multi-nonempty-split*: $M \neq \{\#\} \implies \exists A\ a. M = add\text{-}mset\ a\ A$
 $\langle proof \rangle$

lemma *multiset-cases* [*cases type*]:
obtains (*empty*) $M = \{\#\}$
 $\mid (add)\ x\ N$ **where** $M = add\text{-}mset\ x\ N$
 $\langle proof \rangle$

lemma *multi-drop-mem-not-eq*: $c \in\# B \implies B - \{\#c\} \neq B$
 $\langle proof \rangle$

lemma *multiset-partition*: $M = \{\# x \in\# M. P\ x\ \#\} + \{\# x \in\# M. \neg P\ x\ \#\}$
 $\langle proof \rangle$

lemma *mset-subset-size*: $(A :: 'a\ multiset) \subset\# B \implies size\ A < size\ B$
 $\langle proof \rangle$

lemma *size-1-singleton-mset*: $size\ M = 1 \implies \exists a. M = \{\#a\}$
 $\langle proof \rangle$

1.5.1 Strong induction and subset induction for multisets

Well-foundedness of strict subset relation

lemma *wf-subset-mset-rel*: $wf \{(M, N :: 'a \text{ multiset}). M \subset\# N\}$
 $\langle proof \rangle$

lemma *full-multiset-induct* [*case-names less*]:
assumes *ih*: $\bigwedge B. \forall (A :: 'a \text{ multiset}). A \subset\# B \longrightarrow P A \Longrightarrow P B$
shows $P B$
 $\langle proof \rangle$

lemma *multi-subset-induct* [*consumes 2, case-names empty add*]:
assumes $F \subseteq\# A$
and *empty*: $P \{\#\}$
and *insert*: $\bigwedge a F. a \in\# A \Longrightarrow P F \Longrightarrow P (\text{add-mset } a F)$
shows $P F$
 $\langle proof \rangle$

1.6 The fold combinator

definition *fold-mset* :: $('a \Rightarrow 'b \Rightarrow 'b) \Rightarrow 'b \Rightarrow 'a \text{ multiset} \Rightarrow 'b$
where
 $\text{fold-mset } f s M = \text{Finite-Set.fold } (\lambda x. f x \hat{\wedge} \text{count } M x) s (\text{set-mset } M)$

lemma *fold-mset-empty* [*simp*]: $\text{fold-mset } f s \{\#\} = s$
 $\langle proof \rangle$

context *comp-fun-commute*
begin

lemma *fold-mset-add-mset* [*simp*]: $\text{fold-mset } f s (\text{add-mset } x M) = f x (\text{fold-mset } f s M)$
 $\langle proof \rangle$

corollary *fold-mset-single*: $\text{fold-mset } f s \{\#x\# \} = f x s$
 $\langle proof \rangle$

lemma *fold-mset-fun-left-comm*: $f x (\text{fold-mset } f s M) = \text{fold-mset } f (f x s) M$
 $\langle proof \rangle$

lemma *fold-mset-union* [*simp*]: $\text{fold-mset } f s (M + N) = \text{fold-mset } f (\text{fold-mset } f s M) N$
 $\langle proof \rangle$

lemma *fold-mset-fusion*:
assumes *comp-fun-commute* *g*
and $*$: $\bigwedge x y. h (g x y) = f x (h y)$
shows $h (\text{fold-mset } g w A) = \text{fold-mset } f (h w) A$
 $\langle proof \rangle$

end

lemma *union-fold-mset-add-mset*: $A + B = \text{fold-mset } \text{add-mset } A B$
 ⟨proof⟩

A note on code generation: When defining some function containing a sub-term *fold-mset* F , code generation is not automatic. When interpreting locale *left-commutative* with F , the would be code thms for *fold-mset* become thms like *fold-mset* $F z \{\#\} = z$ where F is not a pattern but contains defined symbols, i.e. is not a code thm. Hence a separate constant with its own code thms needs to be introduced for F . See the image operator below.

1.7 Image

definition *image-mset* :: $('a \Rightarrow 'b) \Rightarrow 'a \text{ multiset} \Rightarrow 'b \text{ multiset}$ **where**
 $\text{image-mset } f = \text{fold-mset } (\text{add-mset} \circ f) \{\#\}$

lemma *comp-fun-commute-mset-image*: $\text{comp-fun-commute } (\text{add-mset} \circ f)$
 ⟨proof⟩

lemma *image-mset-empty* [simp]: $\text{image-mset } f \{\#\} = \{\#\}$
 ⟨proof⟩

lemma *image-mset-single*: $\text{image-mset } f \{\#x\# \} = \{\#f x\# \}$
 ⟨proof⟩

lemma *image-mset-union* [simp]: $\text{image-mset } f (M + N) = \text{image-mset } f M + \text{image-mset } f N$
 ⟨proof⟩

corollary *image-mset-add-mset* [simp]:
 $\text{image-mset } f (\text{add-mset } a M) = \text{add-mset } (f a) (\text{image-mset } f M)$
 ⟨proof⟩

lemma *set-image-mset* [simp]: $\text{set-mset } (\text{image-mset } f M) = \text{image } f (\text{set-mset } M)$
 ⟨proof⟩

lemma *size-image-mset* [simp]: $\text{size } (\text{image-mset } f M) = \text{size } M$
 ⟨proof⟩

lemma *image-mset-is-empty-iff* [simp]: $\text{image-mset } f M = \{\#\} \longleftrightarrow M = \{\#\}$
 ⟨proof⟩

lemma *image-mset-If*:
 $\text{image-mset } (\lambda x. \text{if } P x \text{ then } f x \text{ else } g x) A =$
 $\text{image-mset } f (\text{filter-mset } P A) + \text{image-mset } g (\text{filter-mset } (\lambda x. \neg P x) A)$
 ⟨proof⟩

lemma *image-mset-Diff*:
 assumes $B \subseteq\# A$

shows $image\text{-}mset\ f\ (A - B) = image\text{-}mset\ f\ A - image\text{-}mset\ f\ B$
 ⟨proof⟩

lemma *count-image-mset*:

$count\ (image\text{-}mset\ f\ A)\ x = (\sum y \in f\ -'\ \{x\} \cap set\text{-}mset\ A.\ count\ A\ y)$
 ⟨proof⟩

lemma *image-mset-subseteq-mono*: $A \subseteq\# B \implies image\text{-}mset\ f\ A \subseteq\# image\text{-}mset\ f\ B$

⟨proof⟩

syntax (*ASCII*)

-comprehension-mset :: $'a \Rightarrow 'b \Rightarrow 'b\ multiset \Rightarrow 'a\ multiset\ ((\{\#\text{-}/.\ -:\#\ -\#\}))$

syntax

-comprehension-mset :: $'a \Rightarrow 'b \Rightarrow 'b\ multiset \Rightarrow 'a\ multiset\ ((\{\#\text{-}/.\ -\in\#\ -\#\}))$

translations

$\{\#e.\ x \in\# M\#\} \rightleftharpoons CONST\ image\text{-}mset\ (\lambda x.\ e)\ M$

syntax (*ASCII*)

-comprehension-mset' :: $'a \Rightarrow 'b \Rightarrow 'b\ multiset \Rightarrow bool \Rightarrow 'a\ multiset\ ((\{\#\text{-}/\ |\ -:\#\ -\#\}))$

syntax

-comprehension-mset' :: $'a \Rightarrow 'b \Rightarrow 'b\ multiset \Rightarrow bool \Rightarrow 'a\ multiset\ ((\{\#\text{-}/\ |\ -\in\#\ -\#\}))$

translations

$\{\#e\ |\ x \in\# M.\ P\#\} \rightarrow \{\#e.\ x \in\# \{\#x \in\# M.\ P\#\}\#\}$

This allows to write not just filters like $\{\#x \in\# M.\ x < c\#\}$ but also images like $\{\#x + x.\ x \in\# M\#\}$ and $\{\#x + x\ |\ x \in\# M.\ x < c\#\}$, where the latter is currently displayed as $\{\#x + x.\ x \in\# \{\#x \in\# M.\ x < c\#\}\#\}$.

lemma *in-image-mset*: $y \in\# \{\#f\ x.\ x \in\# M\#\} \longleftrightarrow y \in f\ '\ set\text{-}mset\ M$

⟨proof⟩

functor *image-mset*: *image-mset*

⟨proof⟩

declare

image-mset.id [*simp*]

image-mset.identity [*simp*]

lemma *image-mset-id*[*simp*]: *image-mset id* $x = x$

⟨proof⟩

lemma *image-mset-cong*: $(\bigwedge x.\ x \in\# M \implies f\ x = g\ x) \implies \{\#f\ x.\ x \in\# M\#\} = \{\#g\ x.\ x \in\# M\#\}$

⟨proof⟩

lemma *image-mset-cong-pair*:

$(\forall x\ y.\ (x, y) \in\# M \longrightarrow f\ x\ y = g\ x\ y) \implies \{\#f\ x\ y.\ (x, y) \in\# M\#\} = \{\#g\ x\ y.\ (x, y) \in\# M\#\}$

$y. (x, y) \in \# M \#$
 $\langle \text{proof} \rangle$

1.8 Further conversions

primrec $mset :: 'a \text{ list} \Rightarrow 'a \text{ multiset}$ **where**

$mset [] = \{\#\}$ |
 $mset (a \# x) = \text{add-mset } a (mset x)$

lemma *in-multiset-in-set*:

$x \in \# mset xs \longleftrightarrow x \in \text{set } xs$
 $\langle \text{proof} \rangle$

lemma *count-mset*:

$\text{count } (mset xs) x = \text{length } (\text{filter } (\lambda y. x = y) xs)$
 $\langle \text{proof} \rangle$

lemma *mset-zero-iff[simp]*: $(mset x = \{\#\}) = (x = [])$
 $\langle \text{proof} \rangle$

lemma *mset-zero-iff-right[simp]*: $(\{\#\} = mset x) = (x = [])$
 $\langle \text{proof} \rangle$

lemma *mset-single-iff[iff]*: $mset xs = \{\#x\# \} \longleftrightarrow xs = [x]$
 $\langle \text{proof} \rangle$

lemma *mset-single-iff-right[iff]*: $\{\#x\# \} = mset xs \longleftrightarrow xs = [x]$
 $\langle \text{proof} \rangle$

lemma *set-mset-mset[simp]*: $\text{set-mset } (mset xs) = \text{set } xs$
 $\langle \text{proof} \rangle$

lemma *set-mset-comp-mset [simp]*: $\text{set-mset} \circ mset = \text{set}$
 $\langle \text{proof} \rangle$

lemma *size-mset [simp]*: $\text{size } (mset xs) = \text{length } xs$
 $\langle \text{proof} \rangle$

lemma *mset-append [simp]*: $mset (xs @ ys) = mset xs + mset ys$
 $\langle \text{proof} \rangle$

lemma *mset-filter*: $mset (\text{filter } P xs) = \{\#x \in \# mset xs. P x \#\}$
 $\langle \text{proof} \rangle$

lemma *mset-rev [simp]*:
 $mset (\text{rev } xs) = mset xs$
 $\langle \text{proof} \rangle$

lemma *surj-mset*: *surj* $mset$

<proof>

lemma *distinct-count-atmost-1*:

distinct x = (∀ a. count (mset x) a = (if a ∈ set x then 1 else 0))

<proof>

lemma *mset-eq-setD*:

assumes *mset xs = mset ys*

shows *set xs = set ys*

<proof>

lemma *set-eq-iff-mset-eq-distinct*:

distinct x ⇒ distinct y ⇒

(set x = set y) = (mset x = mset y)

<proof>

lemma *set-eq-iff-mset-remdups-eq*:

(set x = set y) = (mset (remdups x) = mset (remdups y))

<proof>

lemma *mset-compl-union [simp]*: *mset [x←xs. P x] + mset [x←xs. ¬P x] = mset xs*

<proof>

lemma *nth-mem-mset*: *i < length ls ⇒ (ls ! i) ∈# mset ls*

<proof>

lemma *mset-remove1 [simp]*: *mset (remove1 a xs) = mset xs - {#a#}*

<proof>

lemma *mset-eq-length*:

assumes *mset xs = mset ys*

shows *length xs = length ys*

<proof>

lemma *mset-eq-length-filter*:

assumes *mset xs = mset ys*

shows *length (filter (λx. z = x) xs) = length (filter (λy. z = y) ys)*

<proof>

lemma *fold-multiset-equiv*:

assumes *f: ∀x y. x ∈ set xs ⇒ y ∈ set xs ⇒ f x ∘ f y = f y ∘ f x*

and *equiv: mset xs = mset ys*

shows *List.fold f xs = List.fold f ys*

<proof>

lemma *mset-insort [simp]*: *mset (insort x xs) = add-mset x (mset xs)*

<proof>

lemma *mset-map*[simp]: $mset (map f xs) = image-mset f (mset xs)$
 ⟨proof⟩

global-interpretation *mset-set*: folding *add-mset* {#}
defines *mset-set* = folding.F *add-mset* {#}
 ⟨proof⟩

lemma *count-mset-set* [simp]:
 $finite A \implies x \in A \implies count (mset-set A) x = 1$ (is PROP ?P)
 $\neg finite A \implies count (mset-set A) x = 0$ (is PROP ?Q)
 $x \notin A \implies count (mset-set A) x = 0$ (is PROP ?R)
 ⟨proof⟩

lemma *elem-mset-set*[simp, intro]: $finite A \implies x \in\# mset-set A \longleftrightarrow x \in A$
 ⟨proof⟩

lemma *mset-set-Union*:
 $finite A \implies finite B \implies A \cap B = \{\} \implies mset-set (A \cup B) = mset-set A + mset-set B$
 ⟨proof⟩

lemma *filter-mset-mset-set* [simp]:
 $finite A \implies filter-mset P (mset-set A) = mset-set \{x \in A. P x\}$
 ⟨proof⟩

lemma *mset-set-Diff*:
assumes $finite A B \subseteq A$
shows $mset-set (A - B) = mset-set A - mset-set B$
 ⟨proof⟩

lemma *mset-set-set*: $distinct xs \implies mset-set (set xs) = mset xs$
 ⟨proof⟩

context *linorder*
begin

definition *sorted-list-of-multiset* :: 'a multiset \Rightarrow 'a list
where

$sorted-list-of-multiset M = fold-mset insert [] M$

lemma *sorted-list-of-multiset-empty* [simp]:
 $sorted-list-of-multiset \{\# \} = []$
 ⟨proof⟩

lemma *sorted-list-of-multiset-singleton* [simp]:
 $sorted-list-of-multiset \{\#x\# \} = [x]$
 ⟨proof⟩

lemma *sorted-list-of-multiset-insert* [simp]:

sorted-list-of-multiset (*add-mset* x M) = *List.insort* x (*sorted-list-of-multiset* M)
 ⟨*proof*⟩

end

lemma *mset-sorted-list-of-multiset* [*simp*]:
mset (*sorted-list-of-multiset* M) = M
 ⟨*proof*⟩

lemma *sorted-list-of-multiset-mset* [*simp*]:
sorted-list-of-multiset (*mset* xs) = *sort* xs
 ⟨*proof*⟩

lemma *finite-set-mset-mset-set*[*simp*]:
finite $A \implies$ *set-mset* (*mset-set* A) = A
 ⟨*proof*⟩

lemma *mset-set-empty-iff*: *mset-set* $A = \{\#\} \iff A = \{\} \vee$ *infinite* A
 ⟨*proof*⟩

lemma *infinite-set-mset-mset-set*:
 \neg *finite* $A \implies$ *set-mset* (*mset-set* A) = $\{\}$
 ⟨*proof*⟩

lemma *set-sorted-list-of-multiset* [*simp*]:
set (*sorted-list-of-multiset* M) = *set-mset* M
 ⟨*proof*⟩

lemma *sorted-list-of-mset-set* [*simp*]:
sorted-list-of-multiset (*mset-set* A) = *sorted-list-of-set* A
 ⟨*proof*⟩

lemma *mset-upt* [*simp*]: *mset* [$m..<n$] = *mset-set* $\{m..<n\}$
 ⟨*proof*⟩

lemma *image-mset-map-of*:
distinct (*map fst* xs) \implies $\{\#\text{the } (\text{map-of } xs \ i). \ i \in \#\ \text{mset } (\text{map fst } xs)\#\} =$ *mset*
 (*map snd* xs)
 ⟨*proof*⟩

lemma *image-mset-mset-set*:
assumes *inj-on* f A
shows *image-mset* f (*mset-set* A) = *mset-set* ($f \text{ ` } A$)
 ⟨*proof*⟩

1.9 More properties of the replicate and repeat operations

lemma *in-replicate-mset*[*simp*]: $x \in \#\ \text{replicate-mset } n \ y \iff n > 0 \wedge x = y$

<proof>

lemma *set-mset-replicate-mset-subset*[simp]: *set-mset (replicate-mset n x) = (if n = 0 then {} else {x})*
<proof>

lemma *size-replicate-mset*[simp]: *size (replicate-mset n M) = n*
<proof>

lemma *count-le-replicate-mset-subset-eq*: *n ≤ count M x ↔ replicate-mset n x ⊆# M*
<proof>

lemma *filter-eq-replicate-mset*: *{#y ∈# D. y = x#} = replicate-mset (count D x) x*
<proof>

lemma *replicate-count-mset-eq-filter-eq*:
replicate (count (mset xs) k) k = filter (HOL.eq k) xs
<proof>

lemma *replicate-mset-eq-empty-iff* [simp]:
replicate-mset n a = {#} ↔ n = 0
<proof>

lemma *replicate-mset-eq-iff*:
replicate-mset m a = replicate-mset n b ↔
m = 0 ∧ n = 0 ∨ m = n ∧ a = b
<proof>

lemma *repeat-mset-cancel1*: *repeat-mset a A = repeat-mset a B ↔ A = B ∨ a = 0*
<proof>

lemma *repeat-mset-cancel2*: *repeat-mset a A = repeat-mset b A ↔ a = b ∨ A = {#}*
<proof>

lemma *repeat-mset-eq-empty-iff*: *repeat-mset n A = {#} ↔ n = 0 ∨ A = {#}*
<proof>

lemma *image-replicate-mset* [simp]:
image-mset f (replicate-mset n a) = replicate-mset n (f a)
<proof>

1.10 Big operators

locale *comm-monoid-mset = comm-monoid*
begin

interpretation *comp-fun-commute* *f*

<proof>

interpretation *comp?*: *comp-fun-commute* *f* \circ *g*

<proof>

context

begin

definition *F* :: 'a *multiset* \Rightarrow 'a

where *eq-fold*: $F\ M = \text{fold-mset } f\ \mathbf{1}\ M$

lemma *empty* [*simp*]: $F\ \{\#\} = \mathbf{1}$

<proof>

lemma *singleton* [*simp*]: $F\ \{\#x\#} = x$

<proof>

lemma *union* [*simp*]: $F\ (M + N) = F\ M * F\ N$

<proof>

lemma *add-mset* [*simp*]: $F\ (\text{add-mset } x\ N) = x * F\ N$

<proof>

lemma *insert* [*simp*]:

shows $F\ (\text{image-mset } g\ (\text{add-mset } x\ A)) = g\ x * F\ (\text{image-mset } g\ A)$

<proof>

lemma *remove*:

assumes $x \in\# A$

shows $F\ A = x * F\ (A - \{\#x\#})$

<proof>

lemma *neutral*:

$\forall x \in\# A. x = \mathbf{1} \implies F\ A = \mathbf{1}$

<proof>

lemma *neutral-const* [*simp*]:

$F\ (\text{image-mset } (\lambda-. \mathbf{1})\ A) = \mathbf{1}$

<proof> **lemma** *F-image-mset-product*:

$F\ \{\#g\ x\ j * F\ \{\#g\ i\ j. i \in\# A\#\}. j \in\# B\#\} =$

$F\ (\text{image-mset } (g\ x)\ B) * F\ \{\#F\ \{\#g\ i\ j. i \in\# A\#\}. j \in\# B\#\}$

<proof>

lemma *commute*:

$F\ (\text{image-mset } (\lambda i. F\ (\text{image-mset } (g\ i)\ B))\ A) =$

$F\ (\text{image-mset } (\lambda j. F\ (\text{image-mset } (\lambda i. g\ i\ j)\ A))\ B)$

<proof>

lemma *distrib*: $F (\text{image-mset } (\lambda x. g x * h x) A) = F (\text{image-mset } g A) * F (\text{image-mset } h A)$
 ⟨proof⟩

lemma *union-disjoint*:

$A \cap \# B = \{\#\} \implies F (\text{image-mset } g (A \cup \# B)) = F (\text{image-mset } g A) * F (\text{image-mset } g B)$
 ⟨proof⟩

end
end

lemma *comp-fun-commute-plus-mset*[*simp*]: *comp-fun-commute* (*op* + :: 'a multi-set \Rightarrow - \Rightarrow -)
 ⟨proof⟩

declare *comp-fun-commute.fold-mset-add-mset*[*OF comp-fun-commute-plus-mset, simp*]

lemma *in-mset-fold-plus-iff*[*iff*]: $x \in \# \text{fold-mset } (op +) M NN \longleftrightarrow x \in \# M \vee (\exists N. N \in \# NN \wedge x \in \# N)$
 ⟨proof⟩

context *comm-monoid-add*
begin

sublocale *sum-mset*: *comm-monoid-mset plus 0*
defines *sum-mset* = *sum-mset.F* ⟨proof⟩

lemma (**in** *semiring-1*) *sum-mset-replicate-mset* [*simp*]:
sum-mset (*replicate-mset* *n a*) = *of-nat* *n * a*
 ⟨proof⟩

lemma *sum-unfold-sum-mset*:
sum *f A* = *sum-mset* (*image-mset* *f* (*mset-set* *A*))
 ⟨proof⟩

lemma *sum-mset-delta*: *sum-mset* (*image-mset* ($\lambda x. \text{if } x = y \text{ then } c \text{ else } 0$) *A*) = *c * count A y*
 ⟨proof⟩

lemma *sum-mset-delta'*: *sum-mset* (*image-mset* ($\lambda x. \text{if } y = x \text{ then } c \text{ else } 0$) *A*) = *c * count A y*
 ⟨proof⟩

end

lemma *of-nat-sum-mset* [*simp*]:

$of\text{-}nat (sum\text{-}mset M) = sum\text{-}mset (image\text{-}mset of\text{-}nat M)$
 $\langle proof \rangle$

lemma $sum\text{-}mset\text{-}0\text{-}iff$ [simp]:
 $sum\text{-}mset M = (0 :: 'a :: canonically\text{-}ordered\text{-}monoid\text{-}add)$
 $\longleftrightarrow (\forall x \in set\text{-}mset M. x = 0)$
 $\langle proof \rangle$

lemma $sum\text{-}mset\text{-}diff$:
fixes $M N :: ('a :: ordered\text{-}cancel\text{-}comm\text{-}monoid\text{-}diff) multiset$
shows $N \subseteq\# M \implies sum\text{-}mset (M - N) = sum\text{-}mset M - sum\text{-}mset N$
 $\langle proof \rangle$

lemma $size\text{-}eq\text{-}sum\text{-}mset$: $size M = sum\text{-}mset (image\text{-}mset (\lambda\text{-}. 1) M)$
 $\langle proof \rangle$

lemma $size\text{-}mset\text{-}set$ [simp]: $size (mset\text{-}set A) = card A$
 $\langle proof \rangle$

lemma $sum\text{-}mset\text{-}sum\text{-}list$: $sum\text{-}mset (mset xs) = sum\text{-}list xs$
 $\langle proof \rangle$

syntax (ASCII)

$-sum\text{-}mset\text{-}image :: ptrn \Rightarrow 'b set \Rightarrow 'a \Rightarrow 'a :: comm\text{-}monoid\text{-}add ((\exists SUM -\in\# -.$
 $-) [0, 51, 10] 10)$

syntax

$-sum\text{-}mset\text{-}image :: ptrn \Rightarrow 'b set \Rightarrow 'a \Rightarrow 'a :: comm\text{-}monoid\text{-}add ((\exists \sum -\in\# -.$
 $-) [0, 51, 10] 10)$

translations

$\sum i \in\# A. b \Rightarrow CONST sum\text{-}mset (CONST image\text{-}mset (\lambda i. b) A)$

lemma $sum\text{-}mset\text{-}distrib\text{-}left$:

fixes $f :: 'a \Rightarrow 'b :: semiring\text{-}0$
shows $c * (\sum x \in\# M. f x) = (\sum x \in\# M. c * f(x))$
 $\langle proof \rangle$

lemma $sum\text{-}mset\text{-}distrib\text{-}right$:

fixes $f :: 'a \Rightarrow 'b :: semiring\text{-}0$
shows $(\sum b \in\# B. f b) * a = (\sum b \in\# B. f b * a)$
 $\langle proof \rangle$

lemma $sum\text{-}mset\text{-}constant$ [simp]:

fixes $y :: 'b :: semiring\text{-}1$
shows $\langle (\sum x \in\# A. y) = of\text{-}nat (size A) * y \rangle$
 $\langle proof \rangle$

lemma (in $ordered\text{-}comm\text{-}monoid\text{-}add$) $sum\text{-}mset\text{-}mono$:

assumes $\bigwedge i. i \in\# K \implies f i \leq g i$
shows $sum\text{-}mset (image\text{-}mset f K) \leq sum\text{-}mset (image\text{-}mset g K)$

<proof>

lemma *sum-mset-product*:

fixes $f :: 'a::\{\text{comm-monoid-add,times}\} \Rightarrow 'b::\text{semiring-0}$

shows $(\sum i \in\# A. f i) * (\sum i \in\# B. g i) = (\sum i \in\# A. \sum j \in\# B. f i * g j)$

<proof>

abbreviation *Union-mset* :: *'a multiset multiset* \Rightarrow *'a multiset* ($\bigcup\#$ - [900] 900)

where $\bigcup\# MM \equiv \text{sum-mset } MM$ — FIXME ambiguous notation – could likewise refer to $\bigsqcup\#$

lemma *set-mset-Union-mset[simp]*: *set-mset* ($\bigcup\# MM$) = ($\bigcup M \in \text{set-mset } MM. \text{set-mset } M$)

<proof>

lemma *in-Union-mset-iff[iff]*: $x \in\# \bigcup\# MM \longleftrightarrow (\exists M. M \in\# MM \wedge x \in\# M)$

<proof>

lemma *count-sum*:

count (*sum* $f A$) $x = \text{sum } (\lambda a. \text{count } (f a) x) A$

<proof>

lemma *sum-eq-empty-iff*:

assumes *finite* A

shows $\text{sum } f A = \{\#\} \longleftrightarrow (\forall a \in A. f a = \{\#\})$

<proof>

lemma *Union-mset-empty-conv[simp]*: $\bigcup\# M = \{\#\} \longleftrightarrow (\forall i \in\# M. i = \{\#\})$

<proof>

context *comm-monoid-mult*

begin

sublocale *prod-mset: comm-monoid-mset times 1*

defines *prod-mset* = *prod-mset.F* *<proof>*

lemma *prod-mset-empty*:

prod-mset $\{\#\} = 1$

<proof>

lemma *prod-mset-singleton*:

prod-mset $\{\#x\# \} = x$

<proof>

lemma *prod-mset-Un*:

prod-mset $(A + B) = \text{prod-mset } A * \text{prod-mset } B$

<proof>

lemma *prod-mset-replicate-mset* [*simp*]:
 $prod\text{-}mset\ (replicate\text{-}mset\ n\ a) = a \hat{\ } n$
 $\langle proof \rangle$

lemma *prod-unfold-prod-mset*:
 $prod\ f\ A = prod\text{-}mset\ (image\text{-}mset\ f\ (mset\text{-}set\ A))$
 $\langle proof \rangle$

lemma *prod-mset-multiplicity*:
 $prod\text{-}mset\ M = prod\ (\lambda x. x \hat{\ } count\ M\ x)\ (set\text{-}mset\ M)$
 $\langle proof \rangle$

lemma *prod-mset-delta*: $prod\text{-}mset\ (image\text{-}mset\ (\lambda x. if\ x = y\ then\ c\ else\ 1)\ A) =$
 $c \hat{\ } count\ A\ y$
 $\langle proof \rangle$

lemma *prod-mset-delta'*: $prod\text{-}mset\ (image\text{-}mset\ (\lambda x. if\ y = x\ then\ c\ else\ 1)\ A) =$
 $c \hat{\ } count\ A\ y$
 $\langle proof \rangle$

end

syntax (*ASCII*)
 $-prod\text{-}mset\text{-}image :: pttrn \Rightarrow 'b\ set \Rightarrow 'a \Rightarrow 'a::comm\text{-}monoid\text{-}mult\ ((3PROD$
 $-\#-. -) [0, 51, 10] 10)$

syntax
 $-prod\text{-}mset\text{-}image :: pttrn \Rightarrow 'b\ set \Rightarrow 'a \Rightarrow 'a::comm\text{-}monoid\text{-}mult\ ((3\prod -\in\#-.$
 $-) [0, 51, 10] 10)$

translations
 $\prod i \in\# A. b \Leftrightarrow CONST\ prod\text{-}mset\ (CONST\ image\text{-}mset\ (\lambda i. b)\ A)$

lemma (**in** *comm-monoid-mult*) *prod-mset-subset-imp-dvd*:
assumes $A \subseteq\# B$
shows $prod\text{-}mset\ A\ dvd\ prod\text{-}mset\ B$
 $\langle proof \rangle$

lemma (**in** *comm-monoid-mult*) *dvd-prod-mset*:
assumes $x \in\# A$
shows $x\ dvd\ prod\text{-}mset\ A$
 $\langle proof \rangle$

lemma (**in** *semidom*) *prod-mset-zero-iff* [*iff*]:
 $prod\text{-}mset\ A = 0 \iff 0 \in\# A$
 $\langle proof \rangle$

lemma (**in** *semidom-divide*) *prod-mset-diff*:
assumes $B \subseteq\# A$ **and** $0 \notin\# B$
shows $prod\text{-}mset\ (A - B) = prod\text{-}mset\ A\ div\ prod\text{-}mset\ B$
 $\langle proof \rangle$

lemma (in *semidom-divide*) *prod-mset-minus*:
assumes $a \in\# A$ **and** $a \neq 0$
shows $\text{prod-mset } (A - \{\#a\}) = \text{prod-mset } A \text{ div } a$
<proof>

lemma (in *algebraic-semidom*) *is-unit-prod-mset-iff*:
 $\text{is-unit } (\text{prod-mset } A) \longleftrightarrow (\forall x \in\# A. \text{is-unit } x)$
<proof>

lemma (in *normalization-semidom*) *normalize-prod-mset*:
 $\text{normalize } (\text{prod-mset } A) = \text{prod-mset } (\text{image-mset } \text{normalize } A)$
<proof>

lemma (in *normalization-semidom*) *normalized-prod-msetI*:
assumes $\bigwedge a. a \in\# A \implies \text{normalize } a = a$
shows $\text{normalize } (\text{prod-mset } A) = \text{prod-mset } A$
<proof>

lemma *prod-mset-prod-list*: $\text{prod-mset } (\text{mset } xs) = \text{prod-list } xs$
<proof>

1.11 Alternative representations

1.11.1 Lists

context *linorder*
begin

lemma *mset-insort [simp]*:
 $\text{mset } (\text{insort-key } k \ x \ xs) = \text{add-mset } x \ (\text{mset } xs)$
<proof>

lemma *mset-sort [simp]*:
 $\text{mset } (\text{sort-key } k \ xs) = \text{mset } xs$
<proof>

This lemma shows which properties suffice to show that a function f with $f \ xs = f \ ys$ behaves like `sort`.

lemma *properties-for-sort-key*:
assumes $\text{mset } ys = \text{mset } xs$
and $\bigwedge k. k \in \text{set } ys \implies \text{filter } (\lambda x. f \ k = f \ x) \ ys = \text{filter } (\lambda x. f \ k = f \ x) \ xs$
and $\text{sorted } (\text{map } f \ ys)$
shows $\text{sort-key } f \ xs = ys$
<proof>

lemma *properties-for-sort*:
assumes *multiset*: $\text{mset } ys = \text{mset } xs$
and $\text{sorted } ys$
shows $\text{sort } xs = ys$

<proof>

lemma *sort-key-inj-key-eq*:
 assumes *mset-equal*: $mset\ xs = mset\ ys$
 and *inj-on* $f\ (set\ xs)$
 and *sorted* $(map\ f\ ys)$
 shows $sort\text{-}key\ f\ xs = ys$
<proof>

lemma *sort-key-eq-sort-key*:
 assumes $mset\ xs = mset\ ys$
 and *inj-on* $f\ (set\ xs)$
 shows $sort\text{-}key\ f\ xs = sort\text{-}key\ f\ ys$
<proof>

lemma *sort-key-by-quicksort*:
 $sort\text{-}key\ f\ xs = sort\text{-}key\ f\ [x \leftarrow xs.\ f\ x < f\ (xs\ !\ (length\ xs\ div\ 2))]$
 @ $[x \leftarrow xs.\ f\ x = f\ (xs\ !\ (length\ xs\ div\ 2))]$
 @ $sort\text{-}key\ f\ [x \leftarrow xs.\ f\ x > f\ (xs\ !\ (length\ xs\ div\ 2))]$ (**is** $sort\text{-}key\ f\ ?lhs = ?rhs$)
<proof>

lemma *sort-by-quicksort*:
 $sort\ xs = sort\ [x \leftarrow xs.\ x < xs\ !\ (length\ xs\ div\ 2)]$
 @ $[x \leftarrow xs.\ x = xs\ !\ (length\ xs\ div\ 2)]$
 @ $sort\ [x \leftarrow xs.\ x > xs\ !\ (length\ xs\ div\ 2)]$ (**is** $sort\ ?lhs = ?rhs$)
<proof>

A stable parametrized quicksort

definition *part* :: $('b \Rightarrow 'a) \Rightarrow 'a \Rightarrow 'b\ list \Rightarrow 'b\ list \times 'b\ list \times 'b\ list$ **where**
 $part\ f\ pivot\ xs = ([x \leftarrow xs.\ f\ x < pivot], [x \leftarrow xs.\ f\ x = pivot], [x \leftarrow xs.\ pivot < f\ x])$

lemma *part-code* [*code*]:
 $part\ f\ pivot\ [] = ([], [], [])$
 $part\ f\ pivot\ (x \# xs) = (let\ (lts,\ eqs,\ gts) = part\ f\ pivot\ xs;\ x' = f\ x\ in$
 $if\ x' < pivot\ then\ (x \# lts,\ eqs,\ gts)$
 $else\ if\ x' > pivot\ then\ (lts,\ eqs,\ x \# gts)$
 $else\ (lts,\ x \# eqs,\ gts))$
<proof>

lemma *sort-key-by-quicksort-code* [*code*]:
 $sort\text{-}key\ f\ xs =$
 (*case* xs *of*
 $[] \Rightarrow []$
 | $[x] \Rightarrow xs$
 | $[x,\ y] \Rightarrow (if\ f\ x \leq f\ y\ then\ xs\ else\ [y,\ x])$
 | $- \Rightarrow$
 $let\ (lts,\ eqs,\ gts) = part\ f\ (f\ (xs\ !\ (length\ xs\ div\ 2)))\ xs$
 $in\ sort\text{-}key\ f\ lts\ @\ eqs\ @\ sort\text{-}key\ f\ gts$)

<proof>

end

hide-const (open) part

lemma *mset-remdups-subset-eq*: $mset (remdups\ xs) \subseteq\# mset\ xs$
<proof>

lemma *mset-update*:
 $i < length\ ls \implies mset (ls[i := v]) = add-mset\ v (mset\ ls - \{\#ls\ i\#})$
<proof>

lemma *mset-swap*:
 $i < length\ ls \implies j < length\ ls \implies$
 $mset (ls[j := ls\ !\ i, i := ls\ !\ j]) = mset\ ls$
<proof>

1.12 The multiset order

1.12.1 Well-foundedness

definition *mult1* :: $('a \times 'a)$ set $\implies ('a\ multiset \times 'a\ multiset)$ set **where**
 $mult1\ r = \{(N, M). \exists a\ M0\ K. M = add-mset\ a\ M0 \wedge N = M0 + K \wedge$
 $(\forall b. b \in\# K \implies (b, a) \in r)\}$

definition *mult* :: $('a \times 'a)$ set $\implies ('a\ multiset \times 'a\ multiset)$ set **where**
 $mult\ r = (mult1\ r)^+$

lemma *mult1I*:
assumes $M = add-mset\ a\ M0$ **and** $N = M0 + K$ **and** $\bigwedge b. b \in\# K \implies (b, a) \in r$
shows $(N, M) \in mult1\ r$
<proof>

lemma *mult1E*:
assumes $(N, M) \in mult1\ r$
obtains $a\ M0\ K$ **where** $M = add-mset\ a\ M0$ $N = M0 + K$ $\bigwedge b. b \in\# K \implies$
 $(b, a) \in r$
<proof>

lemma *mono-mult1*:
assumes $r \subseteq r'$ **shows** $mult1\ r \subseteq mult1\ r'$
<proof>

lemma *mono-mult*:
assumes $r \subseteq r'$ **shows** $mult\ r \subseteq mult\ r'$
<proof>

lemma *not-less-empty [iff]*: $(M, \{\#\}) \notin mult1\ r$

<proof>

lemma *less-add*:

assumes $mult1: (N, add\text{-}mset\ a\ M0) \in mult1\ r$

shows

$(\exists M. (M, M0) \in mult1\ r \wedge N = add\text{-}mset\ a\ M) \vee$

$(\exists K. (\forall b. b \in\# K \longrightarrow (b, a) \in r) \wedge N = M0 + K)$

<proof>

lemma *all-accessible*:

assumes $wf\ r$

shows $\forall M. M \in Wellfounded.acc\ (mult1\ r)$

<proof>

theorem *wf-mult1*: $wf\ r \implies wf\ (mult1\ r)$

<proof>

theorem *wf-mult*: $wf\ r \implies wf\ (mult\ r)$

<proof>

1.12.2 Closure-free presentation

One direction.

lemma *mult-implies-one-step*:

assumes

trans: $trans\ r$ **and**

MN: $(M, N) \in mult\ r$

shows $\exists I\ J\ K. N = I + J \wedge M = I + K \wedge J \neq \{\#\} \wedge (\forall k \in set\text{-}mset\ K. \exists j \in set\text{-}mset\ J. (k, j) \in r)$

<proof>

lemma *one-step-implies-mult*:

assumes

$J \neq \{\#\}$ **and**

$\forall k \in set\text{-}mset\ K. \exists j \in set\text{-}mset\ J. (k, j) \in r$

shows $(I + K, I + J) \in mult\ r$

<proof>

1.13 The multiset extension is cancellative for multiset union

lemma *mult-cancel*:

assumes *trans* s **and** *irrefl* s

shows $(X + Z, Y + Z) \in mult\ s \longleftrightarrow (X, Y) \in mult\ s$ (**is** $?L \longleftrightarrow ?R$)

<proof>

lemma *mult-cancel-max*:

assumes *trans* s **and** *irrefl* s

shows $(X, Y) \in mult\ s \longleftrightarrow (X - X \cap\# Y, Y - X \cap\# Y) \in mult\ s$ (**is** $?L \longleftrightarrow ?R$)

<proof>

1.14 Quasi-executable version of the multiset extension

Predicate variants of *mult* and the reflexive closure of *mult*, which are executable whenever the given predicate *P* is. Together with the standard code equations for *op* $\cap\#$ and *op* $-$ this should yield quadratic (with respect to calls to *P*) implementations of *multp* and *multeqp*.

definition *multp* :: ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow 'a multiset \Rightarrow 'a multiset \Rightarrow bool **where**
multp *P* *N* *M* =
 (let *Z* = *M* $\cap\#$ *N*; *X* = *M* $-$ *Z* in
 X \neq $\{\#\}$ \wedge (let *Y* = *N* $-$ *Z* in ($\forall y \in \text{set-mset } Y. \exists x \in \text{set-mset } X. P y x$)))

definition *multeqp* :: ('a \Rightarrow 'a \Rightarrow bool) \Rightarrow 'a multiset \Rightarrow 'a multiset \Rightarrow bool **where**
multeqp *P* *N* *M* =
 (let *Z* = *M* $\cap\#$ *N*; *X* = *M* $-$ *Z*; *Y* = *N* $-$ *Z* in
 ($\forall y \in \text{set-mset } Y. \exists x \in \text{set-mset } X. P y x$))

lemma *multp-iff*:

assumes *irrefl* *R* **and** *trans* *R* **and** [*simp*]: $\bigwedge x y. P x y \longleftrightarrow (x, y) \in R$
shows *multp* *P* *N* *M* $\longleftrightarrow (N, M) \in \text{mult } R$ (**is** $?L \longleftrightarrow ?R$)

<proof>

lemma *multeqp-iff*:

assumes *irrefl* *R* **and** *trans* *R* **and** $\bigwedge x y. P x y \longleftrightarrow (x, y) \in R$
shows *multeqp* *P* *N* *M* $\longleftrightarrow (N, M) \in (\text{mult } R)^=$

<proof>

1.14.1 Partial-order properties

lemma (in *preorder*) *mult1-lessE*:

assumes $(N, M) \in \text{mult1 } \{(a, b). a < b\}$
obtains *a* *M0* *K* **where** *M* = *add-mset* *a* *M0* *N* = *M0* + *K*
 a $\notin\#$ *K* \wedge *b* $\in\#$ *K* $\implies b < a$

<proof>

instantiation *multiset* :: (*preorder*) *order*

begin

definition *less-multiset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow bool
where *M'* < *M* $\longleftrightarrow (M', M) \in \text{mult } \{(x', x). x' < x\}$

definition *less-eq-multiset* :: 'a multiset \Rightarrow 'a multiset \Rightarrow bool
where *less-eq-multiset* *M'* *M* $\longleftrightarrow M' < M \vee M' = M$

instance

<proof>

end — FIXME avoid junk stemming from type class interpretation

lemma *mset-le-irrefl* [*elim!*]:
fixes $M :: 'a::preorder\ multiset$
shows $M < M \implies R$
 $\langle proof \rangle$

1.14.2 Monotonicity of multiset union

lemma *mult1-union*: $(B, D) \in mult1\ r \implies (C + B, C + D) \in mult1\ r$
 $\langle proof \rangle$

lemma *union-le-mono2*: $B < D \implies C + B < C + (D::'a::preorder\ multiset)$
 $\langle proof \rangle$

lemma *union-le-mono1*: $B < D \implies B + C < D + (C::'a::preorder\ multiset)$
 $\langle proof \rangle$

lemma *union-less-mono*:
fixes $A\ B\ C\ D :: 'a::preorder\ multiset$
shows $A < C \implies B < D \implies A + B < C + D$
 $\langle proof \rangle$

instantiation *multiset* :: (*preorder*) *ordered-ab-semigroup-add*
begin
instance
 $\langle proof \rangle$
end

1.14.3 Termination proofs with multiset orders

lemma *multi-member-skip*: $x \in\# XS \implies x \in\# \{\# y \#\} + XS$
and *multi-member-this*: $x \in\# \{\# x \#\} + XS$
and *multi-member-last*: $x \in\# \{\# x \#\}$
 $\langle proof \rangle$

definition *ms-strict* = *mult pair-less*
definition *ms-weak* = *ms-strict* \cup *Id*

lemma *ms-reduction-pair*: *reduction-pair* (*ms-strict*, *ms-weak*)
 $\langle proof \rangle$

lemma *smsI*:
 $(set-mset\ A, set-mset\ B) \in max-strict \implies (Z + A, Z + B) \in ms-strict$
 $\langle proof \rangle$

lemma *wmsI*:
 $(set-mset\ A, set-mset\ B) \in max-strict \vee A = \{\#\} \wedge B = \{\#\}$
 $\implies (Z + A, Z + B) \in ms-weak$
 $\langle proof \rangle$

inductive *pw-leq*

where

pw-leq-empty: $pw\text{-leq } \{\#\} \{\#\}$
| *pw-leq-step*: $\llbracket (x,y) \in pair\text{-leq}; pw\text{-leq } X Y \rrbracket \implies pw\text{-leq } (\{\#x\# \} + X) (\{\#y\# \} + Y)$

lemma *pw-leq-lstep*:

$(x, y) \in pair\text{-leq} \implies pw\text{-leq } \{\#x\# \} \{\#y\# \}$
<proof>

lemma *pw-leq-split*:

assumes *pw-leq* $X Y$
shows $\exists A B Z. X = A + Z \wedge Y = B + Z \wedge ((set\text{-mset } A, set\text{-mset } B) \in max\text{-strict} \vee (B = \{\#\} \wedge A = \{\#\}))$
<proof>

lemma

assumes *pwleq*: $pw\text{-leq } Z Z'$
shows *ms-strictI*: $(set\text{-mset } A, set\text{-mset } B) \in max\text{-strict} \implies (Z + A, Z' + B) \in ms\text{-strict}$
and *ms-weakI1*: $(set\text{-mset } A, set\text{-mset } B) \in max\text{-strict} \implies (Z + A, Z' + B) \in ms\text{-weak}$
and *ms-weakI2*: $(Z + \{\#\}, Z' + \{\#\}) \in ms\text{-weak}$
<proof>

lemma *empty-neutral*: $\{\#\} + x = x \ x + \{\#\} = x$
and *nonempty-plus*: $\{\# x \#\} + rs \neq \{\#\}$
and *nonempty-single*: $\{\# x \#\} \neq \{\#\}$
<proof>

<ML>

1.15 Legacy theorem bindings

lemmas *multi-count-eq = multiset-eq-iff* [*symmetric*]

lemma *union-commute*: $M + N = N + (M::'a \text{ multiset})$
<proof>

lemma *union-assoc*: $(M + N) + K = M + (N + (K::'a \text{ multiset}))$
<proof>

lemma *union-lcomm*: $M + (N + K) = N + (M + (K::'a \text{ multiset}))$
<proof>

lemmas *union-ac = union-assoc union-commute union-lcomm add-mset-commute*

lemma *union-right-cancel*: $M + K = N + K \longleftrightarrow M = (N::'a \text{ multiset})$
<proof>

lemma *union-left-cancel*: $K + M = K + N \longleftrightarrow M = (N::'a \text{ multiset})$
<proof>

lemma *multi-union-self-other-eq*: $(A::'a \text{ multiset}) + X = A + Y \implies X = Y$
<proof>

lemma *mset-subset-trans*: $(M::'a \text{ multiset}) \subset\# K \implies K \subset\# N \implies M \subset\# N$
<proof>

lemma *multiset-inter-commute*: $A \cap\# B = B \cap\# A$
<proof>

lemma *multiset-inter-assoc*: $A \cap\# (B \cap\# C) = A \cap\# B \cap\# C$
<proof>

lemma *multiset-inter-left-commute*: $A \cap\# (B \cap\# C) = B \cap\# (A \cap\# C)$
<proof>

lemmas *multiset-inter-ac =*
multiset-inter-commute
multiset-inter-assoc
multiset-inter-left-commute

lemma *mset-le-not-refl*: $\neg M < (M::'a::\text{preorder multiset})$
<proof>

lemma *mset-le-trans*: $K < M \implies M < N \implies K < (N::'a::\text{preorder multiset})$
<proof>

lemma *mset-le-not-sym*: $M < N \implies \neg N < (M::'a::\text{preorder multiset})$
<proof>

lemma *mset-le-asym*: $M < N \implies (\neg P \implies N < (M::'a::\text{preorder multiset})) \implies P$
<proof>

<ML>

1.16 Naive implementation using lists

code-datatype *mset*

lemma [*code*]: $\{\#\} = \text{mset } []$
<proof>

lemma [*code*]: $\text{add-mset } x (\text{mset } xs) = \text{mset } (x \# xs)$
<proof>

lemma [*code*]: $\text{Multiset.is-empty } (\text{mset } xs) \longleftrightarrow \text{List.null } xs$

<proof>

lemma *union-code* [code]: $mset\ xs + mset\ ys = mset\ (xs\ @\ ys)$
<proof>

lemma [code]: $image\ mset\ f\ (mset\ xs) = mset\ (map\ f\ xs)$
<proof>

lemma [code]: $filter\ mset\ f\ (mset\ xs) = mset\ (filter\ f\ xs)$
<proof>

lemma [code]: $mset\ xs - mset\ ys = mset\ (fold\ remove1\ ys\ xs)$
<proof>

lemma [code]:
 $mset\ xs \cap\# mset\ ys =$
 $mset\ (snd\ (fold\ (\lambda x\ (ys,\ zs)).$
 $\text{if } x \in set\ ys \text{ then } (remove1\ x\ ys,\ x \# zs) \text{ else } (ys,\ zs))\ xs\ (ys,\ []))$
<proof>

lemma [code]:
 $mset\ xs \cup\# mset\ ys =$
 $mset\ (case\ prod\ append\ (fold\ (\lambda x\ (ys,\ zs)).\ (remove1\ x\ ys,\ x \# zs))\ xs\ (ys,\ []))$
<proof>

declare *in-multiset-in-set* [code-unfold]

lemma [code]: $count\ (mset\ xs)\ x = fold\ (\lambda y.\ \text{if } x = y \text{ then } Suc \text{ else } id)\ xs\ 0$
<proof>

declare *set-mset-mset* [code]

declare *sorted-list-of-multiset-mset* [code]

lemma [code]: — not very efficient, but representation-ignorant!
 $mset\ set\ A = mset\ (sorted\ list\ of\ set\ A)$
<proof>

declare *size-mset* [code]

fun *subset-eq-mset-impl* :: 'a list \Rightarrow 'a list \Rightarrow bool option **where**
 $subset\ eq\ mset\ impl\ []\ ys = Some\ (ys \neq [])$
 $| subset\ eq\ mset\ impl\ (Cons\ x\ xs)\ ys = (case\ List.extract\ (op = x)\ ys\ of$
 $\quad None \Rightarrow None$
 $\quad | Some\ (ys1,\ -,ys2) \Rightarrow subset\ eq\ mset\ impl\ xs\ (ys1\ @\ ys2))$

lemma *subset-eq-mset-impl*: $(subset\ eq\ mset\ impl\ xs\ ys = None \iff \neg mset\ xs \subseteq\# mset\ ys) \wedge$
 $(subset\ eq\ mset\ impl\ xs\ ys = Some\ True \iff mset\ xs \subset\# mset\ ys) \wedge$

(*subset-eq-mset-impl xs ys = Some False* \longrightarrow *mset xs = mset ys*)
<proof>

lemma [code]: *mset xs $\subseteq\#$ mset ys \longleftrightarrow subset-eq-mset-impl xs ys \neq None*
<proof>

lemma [code]: *mset xs $\subset\#$ mset ys \longleftrightarrow subset-eq-mset-impl xs ys = Some True*
<proof>

instantiation *multiset* :: (*equal*) *equal*
begin

definition

[code del]: *HOL.equal A (B :: 'a multiset) \longleftrightarrow A = B*

lemma [code]: *HOL.equal (mset xs) (mset ys) \longleftrightarrow subset-eq-mset-impl xs ys = Some False*
<proof>

instance
<proof>

end

lemma [code]: *sum-mset (mset xs) = sum-list xs*
<proof>

lemma [code]: *prod-mset (mset xs) = fold times xs 1*
<proof>

Exercise for the casual reader: add implementations for *op \leq* and *op $<$* (multiset order).

Quickcheck generators

definition (in *term-syntax*)

msetify :: 'a::typerep list \times (unit \Rightarrow Code-Evaluation.term)

\Rightarrow 'a multiset \times (unit \Rightarrow Code-Evaluation.term) **where**

[code-unfold]: *msetify xs = Code-Evaluation.valtermify mset {·} xs*

notation *fcomp* (**infixl** $\circ>$ 60)

notation *scomp* (**infixl** $\circ\rightarrow$ 60)

instantiation *multiset* :: (*random*) *random*
begin

definition

Quickcheck-Random.random i = Quickcheck-Random.random i $\circ\rightarrow$ (λ xs. Pair (msetify xs))

instance <proof>

end

no-notation *fcomp* (**infixl** $\circ > 60$)
no-notation *scomp* (**infixl** $\circ \rightarrow 60$)

instantiation *multiset* :: (*full-exhaustive*) *full-exhaustive*
begin

definition *full-exhaustive-multiset* :: ('a multiset \times (unit \Rightarrow term) \Rightarrow (bool \times term list) option) \Rightarrow natural \Rightarrow (bool \times term list) option

where

full-exhaustive-multiset *f* *i* = *Quickcheck-Exhaustive.full-exhaustive* ($\lambda xs. f$ (*msetify* *xs*)) *i*

instance \langle *proof* \rangle

end

hide-const (**open**) *msetify*

1.17 BNF setup

definition *rel-mset* **where**

rel-mset *R* *X* *Y* $\longleftrightarrow (\exists xs\ ys. mset\ xs = X \wedge mset\ ys = Y \wedge list-all2\ R\ xs\ ys)$

lemma *mset-zip-take-Cons-drop-twice*:

assumes *length* *xs* = *length* *ys* *j* \leq *length* *xs*

shows *mset* (*zip* (*take* *j* *xs* @ *x* # *drop* *j* *xs*) (*take* *j* *ys* @ *y* # *drop* *j* *ys*)) =
add-mset (*x*,*y*) (*mset* (*zip* *xs* *ys*))

\langle *proof* \rangle

lemma *ex-mset-zip-left*:

assumes *length* *xs* = *length* *ys* *mset* *xs'* = *mset* *xs*

shows $\exists ys'. length\ ys' = length\ xs' \wedge mset\ (zip\ xs'\ ys') = mset\ (zip\ xs\ ys)$

\langle *proof* \rangle

lemma *list-all2-reorder-left-invariance*:

assumes *rel*: *list-all2* *R* *xs* *ys* **and** *ms-x*: *mset* *xs'* = *mset* *xs*

shows $\exists ys'. list-all2\ R\ xs'\ ys' \wedge mset\ ys' = mset\ ys$

\langle *proof* \rangle

lemma *ex-mset*: $\exists xs. mset\ xs = X$

\langle *proof* \rangle

inductive *pred-mset* :: ('a \Rightarrow bool) \Rightarrow 'a multiset \Rightarrow bool

where

pred-mset *P* {#}

| $\llbracket P\ a; pred-mset\ P\ M \rrbracket \Longrightarrow pred-mset\ P\ (add-mset\ a\ M)$

bnf *'a multiset*
map: image-mset
sets: set-mset
bd: natLeq
wits: {#}
rel: rel-mset
pred: pred-mset
 ⟨*proof*⟩

inductive *rel-mset'*
where

Zero[*intro*]: *rel-mset' R {#} {#}*
 | *Plus*[*intro*]: $\llbracket R \ a \ b; \ rel\text{-}mset' \ R \ M \ N \rrbracket \implies \ rel\text{-}mset' \ R \ (add\text{-}mset \ a \ M) \ (add\text{-}mset \ b \ N)$

lemma *rel-mset-Zero*: *rel-mset R {#} {#}*
 ⟨*proof*⟩

declare *multiset.count*[*simp*]
declare *Abs-multiset-inverse*[*simp*]
declare *multiset.count-inverse*[*simp*]
declare *union-preserves-multiset*[*simp*]

lemma *rel-mset-Plus*:
assumes *ab*: *R a b*
and *MN*: *rel-mset R M N*
shows *rel-mset R (add-mset a M) (add-mset b N)*
 ⟨*proof*⟩

lemma *rel-mset'-imp-rel-mset*: *rel-mset' R M N \implies rel-mset R M N*
 ⟨*proof*⟩

lemma *rel-mset-size*: *rel-mset R M N \implies size M = size N*
 ⟨*proof*⟩

lemma *multiset-induct2*[*case-names empty addL addR*]:
assumes *empty*: *P {#} {#}*
and *addL*: $\bigwedge a \ M \ N. \ P \ M \ N \implies P \ (add\text{-}mset \ a \ M) \ N$
and *addR*: $\bigwedge a \ M \ N. \ P \ M \ N \implies P \ M \ (add\text{-}mset \ a \ N)$
shows *P M N*
 ⟨*proof*⟩

lemma *multiset-induct2-size*[*consumes 1, case-names empty add*]:
assumes *c*: *size M = size N*
and *empty*: *P {#} {#}*
and *add*: $\bigwedge a \ b \ M \ N \ a \ b. \ P \ M \ N \implies P \ (add\text{-}mset \ a \ M) \ (add\text{-}mset \ b \ N)$
shows *P M N*
 ⟨*proof*⟩

lemma *msed-map-invL*:
assumes *image-mset f (add-mset a M) = N*
shows $\exists N1. N = \text{add-mset } (f a) N1 \wedge \text{image-mset } f M = N1$
 $\langle \text{proof} \rangle$

lemma *msed-map-invR*:
assumes *image-mset f M = add-mset b N*
shows $\exists M1 a. M = \text{add-mset } a M1 \wedge f a = b \wedge \text{image-mset } f M1 = N$
 $\langle \text{proof} \rangle$

lemma *msed-rel-invL*:
assumes *rel-mset R (add-mset a M) N*
shows $\exists N1 b. N = \text{add-mset } b N1 \wedge R a b \wedge \text{rel-mset } R M N1$
 $\langle \text{proof} \rangle$

lemma *msed-rel-invR*:
assumes *rel-mset R M (add-mset b N)*
shows $\exists M1 a. M = \text{add-mset } a M1 \wedge R a b \wedge \text{rel-mset } R M1 N$
 $\langle \text{proof} \rangle$

lemma *rel-mset-imp-rel-mset'*:
assumes *rel-mset R M N*
shows *rel-mset' R M N*
 $\langle \text{proof} \rangle$

lemma *rel-mset-rel-mset'*: *rel-mset R M N = rel-mset' R M N*
 $\langle \text{proof} \rangle$

The main end product for *rel-mset*: inductive characterization:

lemmas *rel-mset-induct*[*case-names empty add, induct pred: rel-mset*] =
rel-mset'.induct[*unfolded rel-mset-rel-mset'*[*symmetric*]]

1.18 Size setup

lemma *multiset-size-o-map*:
size-multiset g o image-mset f = size-multiset (g o f)
 $\langle \text{proof} \rangle$

$\langle ML \rangle$

hide-const (**open**) *wcount*

end

2 Additions to Isabelle's Main Theories

theory *Additions-to-Main*
imports $\sim\sim$ */src/HOL/Library/Multiset*

begin

2.1 Addition to Finite-Set Theory

lemma *bound-domain-and-range-impl-finitely-many-functions:*
 finite { $f :: \text{nat} \Rightarrow \text{nat}. (\forall i. f\ i \leq n) \wedge (\forall i \geq m. f\ i = 0)$ }
 <proof>

2.2 Additions to Groups-Big Theory

lemma *sum-card-image:*
 assumes *finite* A
 assumes $\forall s \in A. \forall t \in A. s \neq t \longrightarrow (f\ s) \cap (f\ t) = \{\}$
 shows $\text{sum\ card}\ (f\ `\ A) = \text{sum}\ (\lambda a. \text{card}\ (f\ a))\ A$
 <proof>

lemma *card-Union-image:*
 assumes *finite* S
 assumes $\forall s \in f\ `\ S. \text{finite}\ s$
 assumes $\forall s \in S. \forall t \in S. s \neq t \longrightarrow (f\ s) \cap (f\ t) = \{\}$
 shows $\text{card}\ (\bigcup (f\ `\ S)) = \text{sum}\ (\lambda x. \text{card}\ (f\ x))\ S$
 <proof>

2.3 Addition to Set-Interval Theory

lemma *sum-atMost-remove-nat:*
 assumes $k \leq (n :: \text{nat})$
 shows $(\sum_{i \leq n. f\ i}) = f\ k + (\sum_{i \in \{..n\} - \{k\}. f\ i})$
 <proof>

2.4 Additions to Multiset Theory

lemma *set-mset-Abs-multiset:*
 assumes $f \in \text{multiset}$
 shows $\text{set-mset}\ (\text{Abs-multiset}\ f) = \{x. f\ x > 0\}$
 <proof>

lemma *sum-mset-sum-count:*
 $\text{sum-mset}\ M = (\sum_{i \in \text{set-mset}\ M. \text{count}\ M\ i * i)$
 <proof>

lemma *sum-mset-eq-sum-on-supersets:*
 assumes *finite* A *set-mset* $M \subseteq A$
 shows $(\sum_{i \in \text{set-mset}\ M. \text{count}\ M\ i * i) = (\sum_{i \in A. \text{count}\ M\ i * i)$
 <proof>

end

3 Number Partitions

theory *Number-Partition*
imports *Additions-to-Main*
begin

3.1 Number Partitions as $\text{nat} \Rightarrow \text{nat}$ Functions

definition *partitions* :: $(\text{nat} \Rightarrow \text{nat}) \Rightarrow \text{nat} \Rightarrow \text{bool}$ (**infix** *partitions* 50)
where

$$p \text{ partitions } n = ((\forall i. p \ i \neq 0 \longrightarrow 1 \leq i \wedge i \leq n) \wedge (\sum_{i \leq n}. p \ i * i) = n)$$

lemma *partitionsI*:

assumes $\bigwedge i. p \ i \neq 0 \implies 1 \leq i \wedge i \leq n$

assumes $(\sum_{i \leq n}. p \ i * i) = n$

shows $p \text{ partitions } n$

<proof>

lemma *partitionsE*:

assumes $p \text{ partitions } n$

obtains $\bigwedge i. p \ i \neq 0 \implies 1 \leq i \wedge i \leq n \wedge (\sum_{i \leq n}. p \ i * i) = n$

<proof>

lemma *partitions-zero*:

$p \text{ partitions } 0 \longleftrightarrow p = (\lambda i. 0)$

<proof>

lemma *partitions-one*:

$p \text{ partitions } (\text{Suc } 0) \longleftrightarrow p = (\lambda i. 0)(1 := 1)$

<proof>

3.2 Bounds and Finiteness of Number Partitions

lemma *partitions-imp-finite-elements*:

assumes $p \text{ partitions } n$

shows $\text{finite } \{i. 0 < p \ i\}$

<proof>

lemma *partitions-imp-multiset*:

assumes $p \text{ partitions } n$

shows $p \in \text{multiset}$

<proof>

lemma *partitions-bounds*:

assumes $p \text{ partitions } n$

shows $p \ i \leq n$

<proof>

lemma *partitions-parts-bounded*:

assumes $p \text{ partitions } n$

shows $\text{sum } p \{..n\} \leq n$
<proof>

lemma *finite-partitions*:
finite $\{p. p \text{ partitions } n\}$
<proof>

lemma *finite-partitions-k-parts*:
finite $\{p. p \text{ partitions } n \wedge \text{sum } p \{..n\} = k\}$
<proof>

lemma *partitions-remaining-Max-part*:
assumes $p \text{ partitions } n$
assumes $0 < p \ k$
shows $\forall i. n - k < i \wedge i \neq k \longrightarrow p \ i = 0$
<proof>

3.3 Operations of Number Partitions

lemma *partitions-remove1-bounds*:
assumes $p \text{ partitions } n$
assumes $gr0: 0 < p \ k$
assumes $neq: (p(k := p \ k - 1)) \ i \neq 0$
shows $1 \leq i \wedge i \leq n - k$
<proof>

lemma *partitions-remove1*:
assumes $p \text{ partitions } n$
assumes $gr0: 0 < p \ k$
shows $p(k := p \ k - 1) \text{ partitions } (n - k)$
<proof>

lemma *partitions-insert1*:
assumes $p: p \text{ partitions } n$
assumes $k > 0$
shows $(p(k := p \ k + 1)) \text{ partitions } (n + k)$
<proof>

lemma *count-remove1*:
assumes $p \text{ partitions } n$
assumes $0 < p \ k$
shows $(\sum_{i \leq n - k. (p(k := p \ k - 1)) \ i}) = (\sum_{i \leq n. p \ i}) - 1$
<proof>

lemma *count-insert1*:
assumes $p \text{ partitions } n$
shows $\text{sum } (p(k := p \ k + 1)) \ \{..n + k\} = (\sum_{i \leq n. p \ i}) + 1$
<proof>

lemma *partitions-decrease1*:
 assumes p : p partitions m
 assumes sum : $sum\ p\ \{..m\} = k$
 assumes $p\ 1 = 0$
 shows $(\lambda i. p\ (i + 1))$ partitions $m - k$
 $\langle proof \rangle$

lemma *partitions-increase1*:
 assumes $partitions$: p partitions $m - k$
 assumes k : $sum\ p\ \{..m - k\} = k$
 shows $(\lambda i. p\ (i - 1))$ partitions m
 $\langle proof \rangle$

lemma *count-decrease1*:
 assumes p : p partitions m
 assumes sum : $sum\ p\ \{..m\} = k$
 assumes $p\ 1 = 0$
 shows $sum\ (\lambda i. p\ (i + 1))\ \{..m - k\} = k$
 $\langle proof \rangle$

lemma *count-increase1*:
 assumes $partitions$: p partitions $m - k$
 assumes k : $sum\ p\ \{..m - k\} = k$
 shows $(\sum_{i \leq m}. p\ (i - 1)) = k$
 $\langle proof \rangle$

3.4 Number Partitions as Multisets on Natural Numbers

definition *number-partition* :: $nat \Rightarrow nat\ multiset \Rightarrow bool$
 where
 $number-partition\ n\ N = (sum-mset\ N = n \wedge 0 \notin\# N)$

3.4.1 Relationship to Definition on Functions

lemma *count-partitions-iff*:
 $count\ N\ partitions\ n \iff number-partition\ n\ N$
 $\langle proof \rangle$

lemma *partitions-iff-Abs-multiset*:
 $p\ partitions\ n \iff finite\ \{x. 0 < p\ x\} \wedge number-partition\ n\ (Abs-multiset\ p)$
 $\langle proof \rangle$

lemma *size-nat-multiset-eq*:
 fixes N :: $nat\ multiset$
 assumes $number-partition\ n\ N$
 shows $size\ N = sum\ (count\ N)\ \{..n\}$
 $\langle proof \rangle$

end

4 Euler's Partition Theorem

```
theory Euler-Partition
imports
  Main
  ../Card-Number-Partitions/Number-Partition
begin
```

4.1 Preliminaries

4.1.1 Additions to Divides Theory

```
lemma power-div-nat:
  assumes  $c \leq b$ 
  assumes  $a > 0$ 
  shows  $(a :: nat) ^ b \text{ div } a ^ c = a ^ (b - c)$ 
<proof>
```

4.1.2 Additions to Groups-Big Theory

```
lemma sum-div:
  assumes finite A
  assumes  $\bigwedge a. a \in A \implies (b :: 'b :: semiring-div) \text{ dvd } f a$ 
  shows  $(\sum a \in A. f a) \text{ div } b = (\sum a \in A. (f a) \text{ div } b)$ 
<proof>
```

```
lemma sum-mod:
  assumes finite A
  assumes  $\bigwedge a. a \in A \implies f a \text{ mod } b = (0 :: 'b :: \{semiring-div\})$ 
  shows  $(\sum a \in A. f a) \text{ mod } b = 0$ 
<proof>
```

4.1.3 Additions to Set-Interval Theory

```
lemma geometric-sum-2nat:
   $(\sum i < n. (2 :: nat) ^ i) = (2 ^ n - 1)$ 
<proof>
```

4.1.4 Additions to Nat Theory or Power Theory

```
lemma n-leq-2-pow-n:
   $n \leq 2 ^ n$ 
<proof>
```

4.1.5 Additions to Finite-Set Theory

```
lemma finite-exponents:
  finite  $\{i. 2 ^ i \leq (n :: nat)\}$ 
<proof>
```


4.2 Binary Encoding of Natural Numbers

definition *bitset* :: nat ⇒ nat set

where

$$\text{bitset } n = \{i. \text{ odd } (n \text{ div } (2 \wedge i))\}$$

lemma *in-bitset-bound*:

$$b \in \text{bitset } n \implies 2 \wedge b \leq n$$

⟨proof⟩

lemma *in-bitset-bound-weak*:

$$b \in \text{bitset } n \implies b \leq n$$

⟨proof⟩

lemma *finite-bitset*:

$$\text{finite } (\text{bitset } n)$$

⟨proof⟩

lemma *bitset-0*:

$$\text{bitset } 0 = \{\}$$

⟨proof⟩

lemma *binary-induct* [case-names zero even odd]:

assumes $P (0 :: \text{nat})$

assumes $\bigwedge n. P n \implies P (2 * n)$

assumes $\bigwedge n. P n \implies P (2 * n + 1)$

shows $\bigwedge n. P n$

⟨proof⟩

lemma *bitset-2n*: $\text{bitset } (2 * n) = \text{Suc } \text{' } (\text{bitset } n)$

⟨proof⟩

lemma *bitset-Suc*:

assumes *even* n

shows $\text{bitset } (n + 1) = \text{insert } 0 (\text{bitset } n)$

⟨proof⟩

lemma *bitset-2n1*:

$$\text{bitset } (2 * n + 1) = \text{insert } 0 (\text{Suc } \text{' } (\text{bitset } n))$$

⟨proof⟩

lemma *sum-bitset*:

$$\left(\sum_{i \in \text{bitset } n} 2 \wedge i\right) = n$$

⟨proof⟩

lemma *binarysum-div*:

assumes *finite* B

shows $\left(\sum_{i \in B} (2 :: \text{nat}) \wedge i \text{ div } 2 \wedge j = \left(\sum_{i \in B} \text{if } i < j \text{ then } 0 \text{ else } 2 \wedge (i - j)\right)\right)$

$$\text{(is - = } \left(\sum_{i \in -} ?f i\right))$$

<proof>

lemma *odd-iff*:

assumes *finite B*

shows $\text{odd } (\sum_{i \in B}. \text{if } i < x \text{ then } (0::\text{nat}) \text{ else } 2 \wedge (i - x)) = (x \in B)$ (**is odd**
 $(\sum_{i \in \cdot} ?s \ i) = -$)

<proof>

lemma *bitset-sum*:

assumes *finite B*

shows $\text{bitset } (\sum_{i \in B}. 2 \wedge i) = B$

<proof>

4.3 Decomposition of a Number into a Power of Two and an Odd Number

function (*sequential*) *index* :: *nat* \Rightarrow *nat*

where

index 0 = 0

| *index* n = (*if odd n then 0 else Suc (index (n div 2))*)

<proof>

termination

<proof>

function (*sequential*) *oddp* :: *nat* \Rightarrow *nat*

where

oddp 0 = 0

| *oddp* n = (*if odd n then n else oddp (n div 2)*)

<proof>

termination

<proof>

lemma *odd-oddp*:

$\text{odd } (\text{oddp } n) \longleftrightarrow n \neq 0$

<proof>

lemma *index-oddp-decomposition*:

$n = 2 \wedge (\text{index } n) * \text{oddp } n$

<proof>

lemma *oddp-leq*:

$\text{oddp } n \leq n$

<proof>

lemma *index-oddp-unique*:

assumes *odd (m :: nat) odd m'*

shows $(2 \wedge i * m = 2 \wedge i' * m') \longleftrightarrow (i = i' \wedge m = m')$

<proof>

lemma *index-oddpart*:

assumes *odd m*

shows $\text{index } (2 \wedge i * m) = i \text{ oddpart } (2 \wedge i * m) = m$

<proof>

4.4 Partitions With Only Distinct and Only Odd Parts

definition *odd-of-distinct* :: $(\text{nat} \Rightarrow \text{nat}) \Rightarrow \text{nat} \Rightarrow \text{nat}$

where

$\text{odd-of-distinct } p = (\lambda i. \text{if odd } i \text{ then } (\sum j \mid p (2 \wedge j * i) = 1. 2 \wedge j) \text{ else } 0)$

definition *distinct-of-odd* :: $(\text{nat} \Rightarrow \text{nat}) \Rightarrow \text{nat} \Rightarrow \text{nat}$

where

$\text{distinct-of-odd } p = (\lambda i. \text{if index } i \in \text{bitset } (p (\text{oddpart } i)) \text{ then } 1 \text{ else } 0)$

lemma *odd*:

$\text{odd-of-distinct } p \ i \neq 0 \implies \text{odd } i$

<proof>

lemma *distinct-distinct-of-odd*:

$\text{distinct-of-odd } p \ i \leq 1$

<proof>

lemma *odd-of-distinct*:

assumes $\text{odd-of-distinct } p \ i \neq 0$

assumes $\bigwedge i. p \ i \neq 0 \implies i \leq n$

shows $1 \leq i \wedge i \leq n$

<proof>

lemma *distinct-of-odd*:

assumes $\bigwedge i. p \ i * i \leq n \wedge i. p \ i \neq 0 \implies \text{odd } i$

assumes $\text{distinct-of-odd } p \ i \neq 0$

shows $1 \leq i \wedge i \leq n$

<proof>

lemma *odd-distinct*:

assumes $\bigwedge i. p \ i \neq 0 \implies \text{odd } i$

shows $\text{odd-of-distinct } (\text{distinct-of-odd } p) = p$

<proof>

lemma *distinct-odd*:

assumes $\bigwedge i. p \ i \neq 0 \implies 1 \leq i \wedge i \leq n \wedge i. p \ i \leq 1$

shows $\text{distinct-of-odd } (\text{odd-of-distinct } p) = p$

<proof>

lemma *sum-distinct-of-odd*:

assumes $\bigwedge i. p \ i \neq 0 \implies 1 \leq i \wedge i \leq n$

assumes $\bigwedge i. p\ i * i \leq n$
assumes $\bigwedge i. p\ i \neq 0 \implies \text{odd } i$
shows $(\sum_{i \leq n}. \text{distinct-of-odd } p\ i * i) = (\sum_{i \leq n}. p\ i * i)$
 <proof>

lemma *leq-n*:
assumes $\forall i. 0 < p\ i \longrightarrow 1 \leq i \wedge i \leq (n::\text{nat})$
assumes $(\sum_{i \leq n}. p\ i * i) = n$
shows $p\ i * i \leq n$
 <proof>

lemma *distinct-of-odd-in-distinct-partitions*:
assumes $p \in \{p. p\ \text{partitions } n \wedge (\forall i. p\ i \neq 0 \longrightarrow \text{odd } i)\}$
shows $\text{distinct-of-odd } p \in \{p. p\ \text{partitions } n \wedge (\forall i. p\ i \leq 1)\}$
 <proof>

lemma *odd-of-distinct-in-odd-partitions*:
assumes $p \in \{p. p\ \text{partitions } n \wedge (\forall i. p\ i \leq 1)\}$
shows $\text{odd-of-distinct } p \in \{p. p\ \text{partitions } n \wedge (\forall i. p\ i \neq 0 \longrightarrow \text{odd } i)\}$
 <proof>

4.5 Euler's Partition Theorem

theorem *Euler-partition-theorem*:
 $\text{card } \{p. p\ \text{partitions } n \wedge (\forall i. p\ i \leq 1)\} = \text{card } \{p. p\ \text{partitions } n \wedge (\forall i. p\ i \neq 0 \longrightarrow \text{odd } i)\}$
 (is $\text{card } ?\text{distinct-partitions} = \text{card } ?\text{odd-partitions}$)
 <proof>

end

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