

Dirichlet Series

Manuel Eberl

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Abstract

This entry is a formalisation of much of Chapters 2, 3, and 11 of Apostol's "Introduction to Analytic Number Theory" [1]. This includes:

- Definitions and basic properties for several number-theoretic functions (Euler's φ , Möbius μ , Liouville's λ , the divisor function σ , von Mangoldt's Λ)
- Executable code for most of these functions, the most efficient implementations using the factoring algorithm by Thiemann *et al.*
- Dirichlet products and formal Dirichlet series
- Analytic results connecting convergent formal Dirichlet series to complex functions
- Euler product expansions
- Asymptotic estimates of number-theoretic functions including the density of squarefree integers and the average number of divisors of a natural number

These results are useful as a basis for developing more number-theoretic results, such as the Prime Number Theorem.

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1 Miscellaneous auxiliary facts

theory *Dirichlet-Misc*

imports

HOL-Number-Theory.Number-Theory

begin

lemma

fixes $a\ k :: \text{nat}$

assumes $a > 1\ k > 0$

shows *geometric-sum-nat-aux*: $(a - 1) * (\sum_{i < k}. a^i) = a^k - 1$

and *geometric-sum-nat-dvd*: $a - 1 \text{ dvd } a^k - 1$

and *geometric-sum-nat*: $(\sum_{i < k}. a^i) = (a^k - 1) \text{ div } (a - 1)$

<proof>

lemma *dvd-div-gt0*: $d \text{ dvd } n \implies n > 0 \implies n \text{ div } d > (0 :: \text{nat})$

<proof>

lemma *Set-filter-insert*:

$\text{Set.filter } P (\text{insert } x\ A) = (\text{if } P\ x \text{ then } \text{insert } x (\text{Set.filter } P\ A) \text{ else } \text{Set.filter } P\ A)$

<proof>

lemma *Set-filter-union*: $\text{Set.filter } P (A \cup B) = \text{Set.filter } P\ A \cup \text{Set.filter } P\ B$

<proof>

lemma *Set-filter-empty [simp]*: $\text{Set.filter } P\ \{\} = \{\}$

<proof>

lemma *Set-filter-image*: $\text{Set.filter } P (f\ 'A) = f\ ' \text{Set.filter } (P \circ f)\ A$

<proof>

lemma *Set-filter-cong [cong]*:

$(\bigwedge x. x \in A \implies P\ x \longleftrightarrow Q\ x) \implies A = B \implies \text{Set.filter } P\ A = \text{Set.filter } Q\ B$

<proof>

lemma *inj-on-insert'*: $(\bigwedge B. B \in A \implies x \notin B) \implies \text{inj-on } (\text{insert } x)\ A$

<proof>

lemma

assumes *finite A* $A \neq \{\}$

shows *card-even-subset-aux*: $\text{card } \{B. B \subseteq A \wedge \text{even } (\text{card } B)\} = 2^{\text{card } A - 1}$

and *card-odd-subset-aux*: $\text{card } \{B. B \subseteq A \wedge \text{odd } (\text{card } B)\} = 2^{\text{card } A - 1}$

and *card-even-odd-subset*: $\text{card } \{B. B \subseteq A \wedge \text{even } (\text{card } B)\} = \text{card } \{B. B \subseteq A \wedge \text{odd } (\text{card } B)\}$

<proof>

lemma *bij-betw-prod-divisors-coprime*:
assumes *coprime a (b :: nat)*
shows *bij-betw* $(\lambda x. \text{fst } x * \text{snd } x) (\{d. d \text{ dvd } a\} \times \{d. d \text{ dvd } b\}) \{k. k \text{ dvd } a * b\}$
 $\langle \text{proof} \rangle$

lemma *bij-betw-prime-power-divisors*:
assumes *prime (p :: nat)*
shows *bij-betw* $((\wedge) p) \{..k\} \{d. d \text{ dvd } p \wedge k\}$
 $\langle \text{proof} \rangle$

lemma *sum-divisors-coprime-mult*:
assumes *coprime a (b :: nat)*
shows $(\sum d \mid d \text{ dvd } a * b. f d) = (\sum r \mid r \text{ dvd } a. \sum s \mid s \text{ dvd } b. f (r * s))$
 $\langle \text{proof} \rangle$

end

2 Multiplicative arithmetic functions

theory *Multiplicative-Function*
imports
HOL-Number-Theory.Number-Theory
Dirichlet-Misc
begin

2.1 Definition

locale *multiplicative-function* =
fixes *f :: nat \Rightarrow 'a :: comm-semiring-1*
assumes *zero [simp]: f 0 = 0*
assumes *one [simp]: f 1 = 1*
assumes *mult-coprime-aux: a > 1 \implies b > 1 \implies coprime a b \implies f (a * b) = f a * f b*
begin

lemma *Suc-0 [simp]: f (Suc 0) = 1*
 $\langle \text{proof} \rangle$

lemma *mult-coprime*:
assumes *coprime a b*
shows *f (a * b) = f a * f b*
 $\langle \text{proof} \rangle$

lemma *prod-coprime*:
assumes $\bigwedge x y. x \in A \implies y \in A \implies x \neq y \implies \text{coprime } (g x) (g y)$
shows *f (prod g A) = $(\prod x \in A. f (g x))$*
 $\langle \text{proof} \rangle$

```

lemma prod-prime-factors:
  assumes  $n > 0$ 
  shows  $f\ n = (\prod_{p \in \text{prime-factors } n} f\ (p \wedge \text{multiplicity } p\ n))$ 
  <proof>

lemma multiplicative-sum-divisors: multiplicative-function  $(\lambda n. \sum d \mid d \text{ dvd } n. f\ d)$ 
  <proof>

end

locale multiplicative-function' = multiplicative-function for  $f :: \text{nat} \Rightarrow 'a :: \text{comm-semiring-1}$  +
  fixes f-prime-power ::  $\text{nat} \Rightarrow \text{nat} \Rightarrow 'a$  and f-prime ::  $\text{nat} \Rightarrow 'a$ 
  assumes prime-power:  $\text{prime } p \Longrightarrow k > 0 \Longrightarrow f\ (p \wedge k) = f\text{-prime-power } p\ k$ 
  assumes prime-aux:  $\text{prime } p \Longrightarrow f\text{-prime-power } p\ 1 = f\text{-prime } p$ 
begin

lemma prime:  $\text{prime } p \Longrightarrow f\ p = f\text{-prime } p$ 
  <proof>

lemma prod-prime-factors':
  assumes  $n > 0$ 
  shows  $f\ n = (\prod_{p \in \text{prime-factors } n} f\text{-prime-power } p\ (\text{multiplicity } p\ n))$ 
  <proof>

lemma efficient-code-aux:
  assumes  $n > 0$  set  $ps = (\lambda p. (p, \text{multiplicity } p\ n - 1))$  ' prime-factors  $n$  distinct
   $ps$ 
  shows  $f\ n = (\prod (p,d) \leftarrow ps. f\text{-prime-power } p\ (\text{Suc } d))$ 
  <proof>

lemma efficient-code:
  assumes  $\text{set } (ps\ ()) = (\lambda p. (p, \text{multiplicity } p\ n - 1))$  ' prime-factors  $n$  distinct
   $(ps\ ())$ 
  shows  $f\ n = (\text{if } n = 0 \text{ then } 0 \text{ else } (\prod (p,d) \leftarrow ps\ (). f\text{-prime-power } p\ (\text{Suc } d)))$ 
  <proof>

end

locale completely-multiplicative-function =
  fixes  $f :: \text{nat} \Rightarrow 'a :: \text{comm-semiring-1}$ 
  assumes zero-aux:  $f\ 0 = 0$ 
  assumes one-aux:  $f\ (\text{Suc } 0) = 1$ 
  assumes mult-aux:  $a > 1 \Longrightarrow b > 1 \Longrightarrow f\ (a * b) = f\ a * f\ b$ 
begin

lemma mult:  $f\ (a * b) = f\ a * f\ b$ 

```

<proof>

sublocale *multiplicative-function f*
<proof>

lemma *prod*: $f (\text{prod } g \ A) = (\prod_{x \in A}. f (g \ x))$
<proof>

lemma *power*: $f (n \wedge m) = f \ n \wedge m$
<proof>

lemma *prod-prime-factors'*: $n > 0 \implies f \ n = (\prod_{p \in \text{prime-factors } n}. f \ p \wedge \text{multiplicity } p \ n)$
<proof>

end

locale *completely-multiplicative-function'* =
completely-multiplicative-function f for f :: nat => 'a :: comm-semiring-1 +
fixes *f-prime :: nat => 'a*
assumes *f-prime: prime p => f p = f-prime p*
begin

lemma *prod-prime-factors''*: $n > 0 \implies f \ n = (\prod_{p \in \text{prime-factors } n}. f\text{-prime } p \wedge \text{multiplicity } p \ n)$
<proof>

lemma *efficient-code-aux*:
assumes $n > 0$ *set ps = ($\lambda p. (p, \text{multiplicity } p \ n - 1)$) ' prime-factors n distinct*
ps
shows $f \ n = (\prod (p,d) \leftarrow ps. f\text{-prime } p \wedge \text{Suc } d)$
<proof>

lemma *efficient-code*:
assumes $\text{set } (ps \ ()) = (\lambda p. (p, \text{multiplicity } p \ n - 1))$ *' prime-factors n distinct*
(ps ())
shows $f \ n = (\text{if } n = 0 \text{ then } 0 \text{ else } (\prod (p,d) \leftarrow ps \ (). f\text{-prime } p \wedge \text{Suc } d))$
<proof>

end

lemma *multiplicative-function-eqI*:
assumes *multiplicative-function f multiplicative-function g*
assumes $\bigwedge p \ k. \text{prime } p \implies k > 0 \implies f (p \wedge k) = g (p \wedge k)$
shows $f \ n = g \ n$
<proof>

lemma *multiplicative-function-of-natI*:
multiplicative-function f => multiplicative-function ($\lambda n. \text{of-nat } (f \ n)$)

<proof>

lemma *multiplicative-function-of-natD:*

multiplicative-function ($\lambda n. \text{of-nat } (f \ n) :: 'a :: \{\text{ring-char-0, comm-semiring-1}\}$)

\implies

multiplicative-function f

<proof>

lemma *multiplicative-function-mult:*

assumes *multiplicative-function* f *multiplicative-function* g

shows *multiplicative-function* ($\lambda n. f \ n * g \ n$)

<proof>

lemma *multiplicative-function-inverse:*

fixes $f :: \text{nat} \Rightarrow 'a :: \text{field}$

assumes *multiplicative-function* f

shows *multiplicative-function* ($\lambda n. \text{inverse } (f \ n)$)

<proof>

lemma *multiplicative-function-divide:*

fixes $f :: \text{nat} \Rightarrow 'a :: \text{field}$

assumes *multiplicative-function* f *multiplicative-function* g

shows *multiplicative-function* ($\lambda n. f \ n / g \ n$)

<proof>

lemma *completely-multiplicative-function-mult:*

assumes *completely-multiplicative-function* f *completely-multiplicative-function* g

shows *completely-multiplicative-function* ($\lambda n. f \ n * g \ n$)

<proof>

lemma *completely-multiplicative-function-inverse:*

fixes $f :: \text{nat} \Rightarrow 'a :: \text{field}$

assumes *completely-multiplicative-function* f

shows *completely-multiplicative-function* ($\lambda n. \text{inverse } (f \ n)$)

<proof>

lemma *completely-multiplicative-function-divide:*

fixes $f :: \text{nat} \Rightarrow 'a :: \text{field}$

assumes *completely-multiplicative-function* f *completely-multiplicative-function*

g

shows *completely-multiplicative-function* ($\lambda n. f \ n / g \ n$)

<proof>

lemma (**in** *multiplicative-function*) *completely-multiplicativeI:*

assumes $\bigwedge p \ k. \text{prime } p \implies k > 0 \implies f \ (p \ ^k) = f \ p \ ^k$

shows *completely-multiplicative-function* f

<proof>

2.2 Indicator function

definition $ind :: (nat \Rightarrow bool) \Rightarrow nat \Rightarrow 'a :: semiring-1$ **where**
 $ind\ P\ n = (if\ n > 0 \wedge P\ n\ then\ 1\ else\ 0)$

lemma $ind-0$ [*simp*]: $ind\ P\ 0 = 0$ *<proof>*

lemma $ind-nonzero$: $n > 0 \implies ind\ P\ n = (if\ P\ n\ then\ 1\ else\ 0)$
<proof>

lemma $ind-True$ [*simp*]: $P\ n \implies n > 0 \implies ind\ P\ n = 1$
<proof>

lemma $ind-False$ [*simp*]: $\neg P\ n \implies n > 0 \implies ind\ P\ n = 0$
<proof>

lemma $ind-eq-1-iff$: $ind\ P\ n = 1 \iff n > 0 \wedge P\ n$
<proof>

lemma $ind-eq-0-iff$: $ind\ P\ n = 0 \iff n = 0 \vee \neg P\ n$
<proof>

lemma $multiplicative-function-ind$ [*intro?*]:

assumes $P\ 1 \wedge a\ b. a > 1 \implies b > 1 \implies coprime\ a\ b \implies P\ (a * b) \iff P\ a$
 $\wedge P\ b$

shows $multiplicative-function\ (ind\ P)$
<proof>

end

3 Dirichlet convolution

theory $Dirichlet-Product$

imports

$Complex-Main$

$Multiplicative-Function$

begin

lemma $sum-coprime-dvd-cong$:

$(\sum r \mid r\ dvd\ a. \sum s \mid s\ dvd\ b. f\ r\ s) = (\sum r \mid r\ dvd\ a. \sum s \mid s\ dvd\ b. g\ r\ s)$
if $coprime\ a\ b \wedge r\ s. coprime\ r\ s \implies r\ dvd\ a \implies s\ dvd\ b \implies f\ r\ s = g\ r\ s$
<proof>

definition $dirichlet-prod :: (nat \Rightarrow 'a :: semiring-0) \Rightarrow (nat \Rightarrow 'a) \Rightarrow nat \Rightarrow 'a$
where

$dirichlet-prod\ f\ g = (\lambda n. \sum d \mid d\ dvd\ n. f\ d * g\ (n\ div\ d))$

lemma $sum-divisors-code$:

assumes $n > (0::nat)$

shows $(\sum d \mid d \text{ dvd } n. f d) =$
 $\text{fold-atLeastAtMost-nat } (\lambda d \text{ acc. if } d \text{ dvd } n \text{ then } f d + \text{acc else acc}) 1 n 0$
 ⟨proof⟩

lemma *dirichlet-prod-code* [code]:
 $\text{dirichlet-prod } f g n = (\text{if } n = 0 \text{ then } 0 \text{ else}$
 $\text{fold-atLeastAtMost-nat } (\lambda d \text{ acc. if } d \text{ dvd } n \text{ then } f d * g (n \text{ div } d) + \text{acc else}$
 $\text{acc}) 1 n 0)$
 ⟨proof⟩

lemma *dirichlet-prod-0* [simp]: $\text{dirichlet-prod } f g 0 = 0$
 ⟨proof⟩

lemma *dirichlet-prod-Suc-0* [simp]: $\text{dirichlet-prod } f g (\text{Suc } 0) = f (\text{Suc } 0) * g (\text{Suc } 0)$
 ⟨proof⟩

lemma *dirichlet-prod-cong* [cong]:
assumes $(\bigwedge n. n > 0 \implies f n = f' n) (\bigwedge n. n > 0 \implies g n = g' n)$
shows $\text{dirichlet-prod } f g = \text{dirichlet-prod } f' g'$
 ⟨proof⟩

lemma *dirichlet-prod-altdef1*:
 $\text{dirichlet-prod } f g = (\lambda n. \sum d \mid d \text{ dvd } n. f (n \text{ div } d) * g d)$
 ⟨proof⟩

lemma *dirichlet-prod-altdef2*:
 $\text{dirichlet-prod } f g = (\lambda n. \sum (r,d) \mid r * d = n. f r * g d)$
 ⟨proof⟩

lemma *dirichlet-prod-commutes*:
 $\text{dirichlet-prod } (f :: \text{nat} \Rightarrow 'a :: \text{comm-semiring-0}) g = \text{dirichlet-prod } g f$
 ⟨proof⟩

lemma *finite-divisors-nat'*: $n > (0 :: \text{nat}) \implies \text{finite } \{(a,b). a * b = n\}$
 ⟨proof⟩

lemma *dirichlet-prod-assoc-aux1*:
assumes $n > 0$
shows $\text{dirichlet-prod } f (\text{dirichlet-prod } g h) n =$
 $(\sum (a, b, c) \in \{(a, b, c). a * b * c = n\}. f a * g b * h c)$
 ⟨proof⟩

lemma *dirichlet-prod-assoc-aux2*:
assumes $n > 0$
shows $\text{dirichlet-prod } (\text{dirichlet-prod } f g) h n =$
 $(\sum (a, b, c) \in \{(a, b, c). a * b * c = n\}. f a * g b * h c)$
 ⟨proof⟩

lemma *dirichlet-prod-assoc*:

dirichlet-prod (dirichlet-prod f g) h = dirichlet-prod f (dirichlet-prod g h)
<proof>

lemma *dirichlet-prod-const-right* [*simp*]:

assumes $n > 0$
shows *dirichlet-prod f (λn. if n = Suc 0 then c else 0) n = f n * c*
<proof>

lemma *dirichlet-prod-const-left* [*simp*]:

assumes $n > 0$
shows *dirichlet-prod (λn. if n = Suc 0 then c else 0) g n = c * g n*
<proof>

fun *dirichlet-inverse* :: $(\text{nat} \Rightarrow 'a :: \text{comm-ring-1}) \Rightarrow 'a \Rightarrow \text{nat} \Rightarrow 'a$ **where**

dirichlet-inverse f i n =
(if n = 0 then 0 else if n = 1 then i
*else $-i * (\sum d \mid d \text{ dvd } n \wedge d < n. f (n \text{ div } d) * \text{dirichlet-inverse f i d})$)*

lemma *dirichlet-inverse-induct* [*case-names 0 1 gt1*]:

$P 0 \Longrightarrow P (\text{Suc } 0) \Longrightarrow (\bigwedge n. n > 1 \Longrightarrow (\bigwedge k. k < n \Longrightarrow P k) \Longrightarrow P n) \Longrightarrow P n$
<proof>

lemma *dirichlet-inverse-0* [*simp*]: *dirichlet-inverse f i 0 = 0*

<proof>

lemma *dirichlet-inverse-Suc-0* [*simp*]: *dirichlet-inverse f i (Suc 0) = i*

<proof>

declare *dirichlet-inverse.simps* [*simp del*]

lemma *dirichlet-inverse-gt-1*:

$n > 1 \Longrightarrow \text{dirichlet-inverse f i n} =$
 $-i * (\sum d \mid d \text{ dvd } n \wedge d < n. f (n \text{ div } d) * \text{dirichlet-inverse f i d})$
<proof>

lemma *dirichlet-inverse-cong* [*cong*]:

assumes $\bigwedge n. n > 0 \Longrightarrow f n = f' n \wedge i = i' \wedge n = n'$
shows *dirichlet-inverse f i n = dirichlet-inverse f' i' n'*
<proof>

lemma *dirichlet-inverse-gt-1'*:

assumes $n > 1$
shows *dirichlet-inverse f i n =*
 $-i * \text{dirichlet-prod } (\lambda n. \text{if } n = 1 \text{ then } 0 \text{ else } f n) (\text{dirichlet-inverse f i } n)$
<proof>

lemma *of-int-dirichlet-prod*:

of-int (*dirichlet-prod* *f g n*) = *dirichlet-prod* ($\lambda n.$ *of-int* (*f n*)) ($\lambda n.$ *of-int* (*g n*)) *n*
 ⟨*proof*⟩

lemma *of-int-dirichlet-inverse*:

of-int (*dirichlet-inverse* *f i n*) = *dirichlet-inverse* ($\lambda n.$ *of-int* (*f n*)) (*of-int* *i*) *n*
 ⟨*proof*⟩

lemma *dirichlet-inverse-code* [*code*]:

dirichlet-inverse *f i n* = (*if* *n* = 0 *then* 0 *else if* *n* = 1 *then* *i* *else*
 – *i* * *fold-atLeastAtMost-nat* (λd *acc.* *if* *d* *dvd* *n* *then* *f* (*n* *div* *d*) *
dirichlet-inverse *f i d* + *acc* *else acc*) 1 (*n* – 1) 0)
 ⟨*proof*⟩

lemma *dirichlet-prod-inverse*:

assumes *f* 1 * *i* = 1
shows *dirichlet-prod* *f* (*dirichlet-inverse* *f i*) = ($\lambda n.$ *if* *n* = 1 *then* 1 *else* 0)
 ⟨*proof*⟩

lemma *dirichlet-prod-inverse'*:

assumes *f* 1 * *i* = 1
shows *dirichlet-prod* (*dirichlet-inverse* *f i*) *f* = ($\lambda n.$ *if* *n* = 1 *then* 1 *else* 0)
 ⟨*proof*⟩

lemma *dirichlet-inverse-noninvertible*:

assumes *f* (*Suc* 0) = (0 :: 'a :: {*comm-ring-1*}) *i* = 0
shows *dirichlet-inverse* *f i n* = 0
 ⟨*proof*⟩

lemma *multiplicative-dirichlet-prod*:

assumes *multiplicative-function* *f*
assumes *multiplicative-function* *g*
shows *multiplicative-function* (*dirichlet-prod* *f g*)
 ⟨*proof*⟩

lemma *multiplicative-dirichlet-prodD1*:

fixes *f g* :: *nat* \Rightarrow 'a :: *comm-semiring-1-cancel*
assumes *multiplicative-function* (*dirichlet-prod* *f g*)
assumes *multiplicative-function* *f*
assumes [*simp*]: *g* 0 = 0
shows *multiplicative-function* *g*
 ⟨*proof*⟩

lemma *multiplicative-dirichlet-prodD2*:

fixes *f g* :: *nat* \Rightarrow 'a :: *comm-semiring-1-cancel*
assumes *multiplicative-function* (*dirichlet-prod* *f g*)
assumes *multiplicative-function* *g*
assumes [*simp*]: *f* 0 = 0
shows *multiplicative-function* *f*
 ⟨*proof*⟩

```

lemma multiplicative-dirichlet-inverse:
  assumes multiplicative-function f
  shows multiplicative-function (dirichlet-inverse f 1)
  ⟨proof⟩

lemma dirichlet-prod-prime-power:
  assumes prime p
  shows dirichlet-prod f g (p ^ k) = (∑ i≤k. f (p ^ i) * g (p ^ (k - i)))
  ⟨proof⟩

lemma dirichlet-prod-prime:
  assumes prime p
  shows dirichlet-prod f g p = f 1 * g p + f p * g 1
  ⟨proof⟩

locale multiplicative-dirichlet-prod =
  f: multiplicative-function f + g: multiplicative-function g
  for f g :: nat ⇒ 'a :: comm-semiring-1
begin

  sublocale multiplicative-function dirichlet-prod f g
  ⟨proof⟩

end

locale multiplicative-dirichlet-prod' =
  f: multiplicative-function' f f-prime-power f-prime +
  g: multiplicative-function' g g-prime-power g-prime
  for f g :: nat ⇒ 'a :: comm-semiring-1 and f-prime-power g-prime-power f-prime
  g-prime
begin

  sublocale multiplicative-dirichlet-prod f g ⟨proof⟩

  sublocale multiplicative-function' dirichlet-prod f g
  λp k. f-prime-power p k + g-prime-power p k +
  (∑ i∈{0<..
  λp. f-prime p + g-prime p
  ⟨proof⟩

end

end

```

4 Formal Dirichlet series

```

theory Dirichlet-Series
imports

```

Complex-Main
Dirichlet-Product
Multiplicative-Function
HOL-Computational-Algebra.Computational-Algebra
HOL-Number-Theory.Number-Theory
HOL-Library.FuncSet

begin

A formal Dirichlet series

$$A(s) = \sum_{n=1}^{\infty} \frac{a_n}{n^s}$$

is represented its coefficient sequence starting from 1. For simplicity, we represent this in Isabelle with a function of type $\text{nat} \Rightarrow 'a$ whose value for n is the $n + 1$ -th coefficient.

typedef 'a fds = UNIV :: (nat \Rightarrow 'a) set
 <proof>

setup-lifting type-definition-fds

lift-definition fds-nth :: 'a fds \Rightarrow nat \Rightarrow 'a :: zero is
 $\lambda f :: \text{nat} \Rightarrow 'a. \text{case-nat } 0 f$ <proof>

lift-definition fds :: (nat \Rightarrow 'a) \Rightarrow 'a fds is
 $\lambda f. f \circ \text{Suc}$ <proof>

lemma fds-nth-fds: $\text{fds-nth } (f \text{ds } f) n = (\text{if } n = 0 \text{ then } 0 \text{ else } f n)$
 <proof>

lemma fds-nth-fds': $f 0 = 0 \implies \text{fds-nth } (f \text{ds } f) = f$
 <proof>

lemma fds-nth-0 [simp]: $\text{fds-nth } f 0 = 0$
 <proof>

lemma fds-nth-fds-pos [simp]: $n > 0 \implies \text{fds-nth } (f \text{ds } f) n = f n$
 <proof>

lemma fds-fds-nth [simp]: $\text{fds } (\text{fds-nth } f) = f$
 <proof>

lemma fds-eq-fds-iff:
 $\text{fds } f = \text{fds } g \iff (\forall n > 0. f n = g n)$
 <proof>

lemma fds-eq-fds-iff': $f 0 = g 0 \implies \text{fds } f = \text{fds } g \iff f = g$
 <proof>

lemma fds-eqI [intro?]:

assumes $(\bigwedge n. n > 0 \implies \text{fds-nth } f \ n = \text{fds-nth } g \ n)$
shows $f = g$
 $\langle \text{proof} \rangle$

lemma *fds-cong* [*cong*]: $(\bigwedge n. n > 0 \implies f \ n = (g \ n :: 'a :: \text{zero})) \implies \text{fds } f = \text{fds } g$
 $\langle \text{proof} \rangle$

lemma *fds-eq-iff*: $f = g \iff (\forall n > 0. \text{fds-nth } f \ n = \text{fds-nth } g \ n)$
 $\langle \text{proof} \rangle$

lemma *dirichlet-prod-fds-nth-fds-left* [*simp*]:
 $\text{dirichlet-prod } (\text{fds-nth } (\text{fds } f)) \ g = \text{dirichlet-prod } f \ g$
 $\langle \text{proof} \rangle$

lemma *dirichlet-prod-fds-nth-fds-right* [*simp*]:
 $\text{dirichlet-prod } f \ (\text{fds-nth } (\text{fds } g)) = \text{dirichlet-prod } f \ g$
 $\langle \text{proof} \rangle$

definition *fds-const* :: $'a :: \text{zero} \implies 'a \ \text{fds}$ **where**
 $\text{fds-const } c = \text{fds } (\lambda n. \text{if } n = 1 \text{ then } c \ \text{else } 0)$

abbreviation *fds-ind* **where** $\text{fds-ind } P \equiv \text{fds } (\text{ind } P)$

bundle *fds-syntax*
begin

notation *fds-nth* (**infixl** $\langle \$ \rangle$ 75)
notation *fds* (**binder** $\langle \chi \rangle$ 10)
notation *dirichlet-prod* (**infixl** $\langle * \rangle$ 70)

end

instantiation *fds* :: $(\text{zero}) \ \text{zero}$
begin
definition *zero-fds* :: $'a \ \text{fds}$ **where** $\text{zero-fds} = \text{fds } (\lambda -. 0)$
instance $\langle \text{proof} \rangle$
end

instantiation *fds* :: $(\{\text{zero}, \text{one}\}) \ \text{one}$
begin
definition *one-fds* :: $'a \ \text{fds}$ **where** $\text{one-fds} = \text{fds } (\lambda n. \text{if } n = 1 \text{ then } 1 \ \text{else } 0)$
instance $\langle \text{proof} \rangle$
end

instantiation *fds* :: $(\{\text{plus}, \text{zero}\}) \ \text{plus}$
begin

definition *plus-fds* :: 'a fds \Rightarrow 'a fds \Rightarrow 'a fds
where *plus-fds* f g = fds ($\lambda n.$ fds-nth f n + fds-nth g n)
instance <proof>
end

instantiation *fds* :: (semiring-0) times
begin
definition *times-fds* :: 'a fds \Rightarrow 'a fds \Rightarrow 'a fds
where *times-fds* f g = fds (dirichlet-prod (fds-nth f) (fds-nth g))
instance <proof>
end

instantiation *fds* :: ({uminus,zero}) uminus
begin
definition *uminus-fds* :: 'a fds \Rightarrow 'a fds
where *uminus-fds* f = fds ($\lambda n.$ -fds-nth f n)
instance <proof>
end

instantiation *fds* :: ({minus,zero}) minus
begin
definition *minus-fds* :: 'a fds \Rightarrow 'a fds \Rightarrow 'a fds
where *minus-fds* f g = fds ($\lambda n.$ fds-nth f n - fds-nth g n)
instance <proof>
end

4.1 General properties

lemma *fds-nth-zero* [simp]: *fds-nth* 0 = ($\lambda.$ 0)
<proof>

lemma *fds-nth-one*: *fds-nth* 1 = ($\lambda n.$ if n = 1 then 1 else 0)
<proof>

lemma *fds-nth-one-Suc-0* [simp]: *fds-nth* 1 (Suc 0) = 1
<proof>

lemma *fds-nth-one-not-Suc-0* [simp]: $n \neq \text{Suc } 0 \implies \text{fds-nth } 1 \ n = 0$
<proof>

lemma *fds-nth-plus* [simp]:
fds-nth (f + g) = ($\lambda n.$ *fds-nth* f n + *fds-nth* g n :: 'a :: monoid-add)
<proof>

lemma *fds-nth-minus* [simp]:
fds-nth (f - g) = ($\lambda n.$ *fds-nth* f n - *fds-nth* g n :: 'a :: {cancel-comm-monoid-add})
<proof>

lemma *fds-nth-uminus* [simp]: *fds-nth* (-g) = ($\lambda n.$ - *fds-nth* g n :: 'a :: group-add)

<proof>

lemma *fds-nth-mult*: $fds\text{-}nth\ (f * g) = dirichlet\text{-}prod\ (fds\text{-}nth\ f)\ (fds\text{-}nth\ g)$
<proof>

lemma *fds-nth-mult-const-left* [*simp*]: $fds\text{-}nth\ (fds\text{-}const\ c * f)\ n = c * fds\text{-}nth\ f\ n$
<proof>

lemma *fds-nth-mult-const-right* [*simp*]: $fds\text{-}nth\ (f * fds\text{-}const\ c)\ n = fds\text{-}nth\ f\ n * c$
<proof>

instance *fds* :: (*{semigroup-add, zero}*) *semigroup-add*
<proof>

instance *fds* :: (*{ab-semigroup-add, zero}*) *ab-semigroup-add*
<proof>

instance *fds* :: (*{cancel-semigroup-add, zero}*) *cancel-semigroup-add*
<proof>

instance *fds* :: (*{cancel-ab-semigroup-add, zero}*) *cancel-ab-semigroup-add*
<proof>

instance *fds* :: (*monoid-add*) *monoid-add*
<proof>

instance *fds* :: (*comm-monoid-add*) *comm-monoid-add*
<proof>

instance *fds* :: (*cancel-comm-monoid-add*) *cancel-comm-monoid-add*
<proof>

instance *fds* :: (*group-add*) *group-add*
<proof>

instance *fds* :: (*ab-group-add*) *ab-group-add*
<proof>

instance *fds* :: (*semiring-0*) *semiring-0*
<proof>

instance *fds* :: (*comm-semiring-0*) *comm-semiring-0*
<proof>

instance *fds* :: (*semiring-0-cancel*) *semiring-0-cancel*
<proof>

instance *fds* :: (*comm-semiring-0-cancel*) *comm-semiring-0-cancel* \langle *proof* \rangle

instance *fds* :: (*semiring-1*) *semiring-1* \langle *proof* \rangle

instance *fds* :: (*comm-semiring-1*) *comm-semiring-1* \langle *proof* \rangle

instance *fds* :: (*semiring-1-cancel*) *semiring-1-cancel* \langle *proof* \rangle

instance *fds* :: (*ring*) *ring* \langle *proof* \rangle

instance *fds* :: (*ring-1*) *ring-1* \langle *proof* \rangle

instance *fds* :: (*comm-ring*) *comm-ring* \langle *proof* \rangle

instance *fds* :: (*semiring-no-zero-divisors*) *semiring-no-zero-divisors* \langle *proof* \rangle

instance *fds* :: (*ring-no-zero-divisors*) *ring-no-zero-divisors* \langle *proof* \rangle

instance *fds* :: (*idom*) *idom* \langle *proof* \rangle

instantiation *fds* :: (*real-vector*) *real-vector*

begin

definition *scaleR-fds* :: *real* \Rightarrow 'a *fds* \Rightarrow 'a *fds* **where**
scaleR-fds *c* *f* = *fds* ($\lambda n.$ *c* *_R *fds-nth* *f* *n*)

lemma *fds-nth-scaleR* [*simp*]: *fds-nth* (*c* *_R *f*) = ($\lambda n.$ *c* *_R *fds-nth* *f* *n*) \langle *proof* \rangle

instance \langle *proof* \rangle

end

instance *fds* :: (*real-algebra*) *real-algebra* \langle *proof* \rangle

instance *fds* :: (*real-algebra-1*) *real-algebra-1* \langle *proof* \rangle

lemma *fds-nth-sum* [*simp*]: *fds-nth* (*sum* *f* *A*) *n* = *sum* ($\lambda x.$ *fds-nth* (*f* *x*) *n*) *A* \langle *proof* \rangle

lemma *sum-fds* [*simp*]: ($\sum x \in A.$ *fds* (*f* *x*)) = *fds* ($\lambda n.$ $\sum x \in A.$ *f* *x* *n*) \langle *proof* \rangle

lemma *fds-nth-const*: *fds-nth* (*fds-const* *c*) = ($\lambda n.$ *if* *n* = 1 *then* *c* *else* 0) \langle *proof* \rangle

lemma *fds-nth-const-Suc-0* [*simp*]: *fds-nth* (*fds-const* *c*) (*Suc* 0) = *c*

<proof>

lemma *fds-nth-const-not-Suc-0* [*simp*]: $n \neq 1 \implies \text{fds-nth } (\text{fds-const } c) n = 0$
<proof>

lemma *fds-const-zero* [*simp*]: $\text{fds-const } 0 = 0$
<proof>

lemma *fds-const-one* [*simp*]: $\text{fds-const } 1 = 1$
<proof>

lemma *fds-const-add* [*simp*]: $\text{fds-const } (a + b :: 'a :: \text{monoid-add}) = \text{fds-const } a + \text{fds-const } b$
<proof>

lemma *fds-const-minus* [*simp*]:
 $\text{fds-const } (a - b :: 'a :: \text{cancel-comm-monoid-add}) = \text{fds-const } a - \text{fds-const } b$
<proof>

lemma *fds-const-uminus* [*simp*]:
 $\text{fds-const } (- b :: 'a :: \text{ab-group-add}) = - \text{fds-const } b$
<proof>

lemma *fds-const-mult* [*simp*]:
 $\text{fds-const } (a * b :: 'a :: \text{semiring-0}) = \text{fds-const } a * \text{fds-const } b$
<proof>

lemma *fds-const-of-nat* [*simp*]: $\text{fds-const } (\text{of-nat } c) = \text{of-nat } c$
<proof>

lemma *fds-const-of-int* [*simp*]: $\text{fds-const } (\text{of-int } c) = \text{of-int } c$
<proof>

lemma *fds-const-of-real* [*simp*]: $\text{fds-const } (\text{of-real } c) = \text{of-real } c$
<proof>

instantiation *fds* :: ($\{\text{inverse, comm-ring-1}\}$) *inverse*
begin

definition *inverse-fds* :: $'a \text{ fds} \Rightarrow 'a \text{ fds}$ **where**
 $\text{inverse-fds } f = \text{fds } (\lambda n. \text{dirichlet-inverse } (\text{fds-nth } f) (\text{inverse } (\text{fds-nth } f 1)) n)$

definition *divide-fds* :: $'a \text{ fds} \Rightarrow 'a \text{ fds} \Rightarrow 'a \text{ fds}$ **where**
 $\text{divide-fds } f g = f * \text{inverse } g$

instance *<proof>*

end

lemma *numeral-fds*: $\text{numeral } n = \text{fds-const } (\text{numeral } n)$
<proof>

lemma *fds-ind-False* [simp]: $\text{fds-ind } (\lambda-. \text{False}) = 0$
<proof>

lemma *fds-commutes*:

assumes $\bigwedge m n. m > 0 \implies n > 0 \implies \text{fds-nth } f m * \text{fds-nth } g n = \text{fds-nth } g n$
 $* \text{fds-nth } f m$
shows $f * g = g * f$
<proof>

lemma *fds-nth-mult-Suc-0* [simp]:
 $\text{fds-nth } (f * g) (\text{Suc } 0) = \text{fds-nth } f (\text{Suc } 0) * \text{fds-nth } g (\text{Suc } 0)$
<proof>

lemma *fds-nth-inverse*:
 $\text{fds-nth } (\text{inverse } f) = \text{dirichlet-inverse } (\text{fds-nth } f) (\text{inverse } (\text{fds-nth } f 1))$
<proof>

lemma *inverse-fds-nonunit*:
 $\text{fds-nth } f 1 = (0 :: 'a :: \text{field}) \implies \text{inverse } f = 0$
<proof>

lemma *inverse-0-fds* [simp]: $\text{inverse } (0 :: 'a :: \text{field } \text{fds}) = 0$
<proof>

lemma *fds-left-inverse*:
 $\text{fds-nth } f 1 \neq (0 :: 'a :: \text{field}) \implies \text{inverse } f * f = 1$
<proof>

lemma *fds-right-inverse*:
 $\text{fds-nth } f 1 \neq (0 :: 'a :: \text{field}) \implies f * \text{inverse } f = 1$
<proof>

lemma *fds-left-inverse-unique*:
assumes $f * g = (1 :: 'a :: \text{field } \text{fds})$
shows $f = \text{inverse } g$
<proof>

lemma *fds-right-inverse-unique*:
assumes $f * g = (1 :: 'a :: \text{field } \text{fds})$
shows $g = \text{inverse } f$
<proof>

lemma *inverse-1-fds* [simp]: $\text{inverse } (1 :: 'a :: \text{field } \text{fds}) = 1$
<proof>

lemma *inverse-const-fds* [simp]:
 $inverse (fds\text{-}const\ c :: 'a :: field\ fds) = fds\text{-}const\ (inverse\ c)$
 ⟨proof⟩

lemma *inverse-mult-fds*: $inverse\ (f * g :: 'a :: field\ fds) = inverse\ f * inverse\ g$
 ⟨proof⟩

definition *fds-zeta* :: 'a :: one fds
 where $fds\text{-}zeta = fds\ (\lambda\cdot.\ 1)$

lemma *fds-zeta-altdef*: $fds\text{-}zeta = fds\ (\lambda n.\ if\ n = 0\ then\ 0\ else\ 1)$
 ⟨proof⟩

lemma *fds-nth-zeta*: $fds\text{-}nth\ fds\text{-}zeta = (\lambda n.\ if\ n = 0\ then\ 0\ else\ 1)$
 ⟨proof⟩

lemma *fds-nth-zeta-pos* [simp]: $n > 0 \implies fds\text{-}nth\ fds\text{-}zeta\ n = 1$
 ⟨proof⟩

lemma *fds-zeta-commutes*: $fds\text{-}zeta * (f :: 'a :: semiring-1\ fds) = f * fds\text{-}zeta$
 ⟨proof⟩

lemma *fds-ind-True* [simp]: $fds\text{-}ind\ (\lambda\cdot.\ True) = fds\text{-}zeta$
 ⟨proof⟩

lemma *finite-extensional-prod-nat*:
 assumes $finite\ A\ b > 0$
 shows $finite\ \{d \in extensional\ A.\ prod\ d\ A = (b :: nat)\}$
 ⟨proof⟩

The n -th coefficient of a product of Dirichlet series can be determined by summing over all products of k_i -th coefficients of the series such that the product of the k_i is n .

lemma *fds-nth-prod*:
 assumes $finite\ A\ A \neq \{\}\ n > 0$
 shows $fds\text{-}nth\ (\prod_{x \in A} f\ x)\ n =$
 $(\sum d \mid d \in extensional\ A \wedge prod\ d\ A = n.\ \prod_{x \in A} fds\text{-}nth\ (f\ x)\ (d\ x))$
 ⟨proof⟩

lemma *fds-nth-power-Suc-0* [simp]: $fds\text{-}nth\ (f \wedge^n)\ (Suc\ 0) = fds\text{-}nth\ f\ (Suc\ 0) \wedge^n$
 ⟨proof⟩

lemma *fds-nth-prod-Suc-0* [simp]: $fds\text{-}nth\ (prod\ f\ A)\ (Suc\ 0) = (\prod_{x \in A} fds\text{-}nth\ (f\ x)\ (Suc\ 0))$
 ⟨proof⟩

lemma *fds-nth-power-eq-0*:

assumes $n < 2 \wedge k \text{ fds-nth } f \ 1 = 0$
shows $\text{fds-nth } (f \wedge k) \ n = 0$
 $\langle \text{proof} \rangle$

4.2 Shifting the argument

class *nat-power* = *semiring-1* +
fixes *nat-power* :: $\text{nat} \Rightarrow 'a \Rightarrow 'a$
assumes *nat-power-0-left* [*simp*]: $x \neq 0 \implies \text{nat-power } 0 \ x = 0$
assumes *nat-power-0-right* [*simp*]: $n > 0 \implies \text{nat-power } n \ 0 = 1$
assumes *nat-power-1-left* [*simp*]: $\text{nat-power } (\text{Suc } 0) \ x = 1$
assumes *nat-power-1-right* [*simp*]: $\text{nat-power } n \ 1 = \text{of-nat } n$
assumes *nat-power-add*: $n > 0 \implies \text{nat-power } n \ (a + b) = \text{nat-power } n \ a * \text{nat-power } n \ b$
assumes *nat-power-mult-distrib*:
 $m > 0 \implies n > 0 \implies \text{nat-power } (m * n) \ a = \text{nat-power } m \ a * \text{nat-power } n \ a$
assumes *nat-power-power*:
 $n > 0 \implies \text{nat-power } n \ (a * \text{of-nat } m) = \text{nat-power } n \ a \wedge m$
begin

lemma *nat-power-of-nat* [*simp*]: $m > 0 \implies \text{nat-power } m \ (\text{of-nat } n) = \text{of-nat } (m \wedge n)$
 $\langle \text{proof} \rangle$

lemma *nat-power-power-left*: $m > 0 \implies \text{nat-power } (m \wedge k) \ n = \text{nat-power } m \ n \wedge k$
 $\langle \text{proof} \rangle$

end

class *nat-power-field* = *nat-power* + *field* +
assumes *nat-power-nonzero* [*simp*]: $n > 0 \implies \text{nat-power } n \ z \neq 0$
begin

lemma *nat-power-diff*: $n > 0 \implies \text{nat-power } n \ (a - b) = \text{nat-power } n \ a / \text{nat-power } n \ b$
 $\langle \text{proof} \rangle$

end

instantiation *nat* :: *nat-power*

begin

definition [*simp*]: $\text{nat-power-nat } a \ b = (a \wedge b \ :: \ \text{nat})$

instance $\langle \text{proof} \rangle$

end

instantiation *real* :: *nat-power-field*

begin

definition [*simp*]: $\text{nat-power-real } a \ b = (\text{real } a \ \text{powr } b)$

instance $\langle proof \rangle$
end

The following operation corresponds to shifting the argument of a Dirichlet series, i. e. subtracting a constant from it. In effect, this turns the series

$$A(s) = \sum_{n=1}^{\infty} \frac{a_n}{n^s}$$

into the series

$$A(s - c) = \sum_{n=1}^{\infty} \frac{n^c \cdot a_n}{n^s} .$$

definition $fds\text{-}shift :: 'a :: nat\text{-}power \Rightarrow 'a\ fds \Rightarrow 'a\ fds$ **where**
 $fds\text{-}shift\ c\ f = fds\ (\lambda n. fds\text{-}nth\ f\ n * nat\text{-}power\ n\ c)$

lemma $fds\text{-}nth\text{-}shift$ $[simp]$: $fds\text{-}nth\ (fds\text{-}shift\ c\ f)\ n = fds\text{-}nth\ f\ n * nat\text{-}power\ n\ c$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}shift$ $[simp]$: $fds\text{-}shift\ c\ (fds\text{-}shift\ c'\ f) = fds\text{-}shift\ (c' + c)\ f$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}zero$ $[simp]$: $fds\text{-}shift\ c\ 0 = 0$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}1$ $[simp]$: $fds\text{-}shift\ a\ 1 = 1$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}const$ $[simp]$: $fds\text{-}shift\ a\ (fds\text{-}const\ c) = fds\text{-}const\ c$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}add$ $[simp]$:
fixes $f\ g :: 'a :: \{monoid\text{-}add, nat\text{-}power\}$ fds
shows $fds\text{-}shift\ c\ (f + g) = fds\text{-}shift\ c\ f + fds\text{-}shift\ c\ g$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}minus$ $[simp]$:
fixes $f\ g :: 'a :: \{comm\text{-}semiring\text{-}1\text{-}cancel, nat\text{-}power\}$ fds
shows $fds\text{-}shift\ c\ (f - g) = fds\text{-}shift\ c\ f - fds\text{-}shift\ c\ g$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}uminus$ $[simp]$:
fixes $f :: 'a :: \{ring, nat\text{-}power\}$ fds
shows $fds\text{-}shift\ c\ (-f) = -fds\text{-}shift\ c\ f$
 $\langle proof \rangle$

lemma $fds\text{-}shift\text{-}mult$ $[simp]$:
fixes $f\ g :: 'a :: \{comm\text{-}semiring, nat\text{-}power\}$ fds
shows $fds\text{-}shift\ c\ (f * g) = fds\text{-}shift\ c\ f * fds\text{-}shift\ c\ g$

<proof>

lemma *fds-shift-power* [*simp*]:
fixes $f :: 'a :: \{\text{comm-semiring, nat-power}\}$ *fds*
shows $\text{fds-shift } c (f \wedge n) = \text{fds-shift } c f \wedge n$
<proof>

lemma *fds-shift-by-0* [*simp*]: $\text{fds-shift } 0 f = f$
<proof>

lemma *fds-shift-inverse* [*simp*]:
 $\text{fds-shift } (a :: 'a :: \{\text{field, nat-power}\}) (\text{inverse } f) = \text{inverse } (\text{fds-shift } a f)$
<proof>

lemma *fds-shift-divide* [*simp*]:
 $\text{fds-shift } (a :: 'a :: \{\text{field, nat-power}\}) (f / g) = \text{fds-shift } a f / \text{fds-shift } a g$
<proof>

lemma *fds-shift-sum* [*simp*]: $\text{fds-shift } a (\sum x \in A. f x) = (\sum x \in A. \text{fds-shift } a (f x))$
<proof>

lemma *fds-shift-prod* [*simp*]: $\text{fds-shift } a (\prod x \in A. f x) = (\prod x \in A. \text{fds-shift } a (f x))$
<proof>

4.3 Scaling the argument

The following operation corresponds to scaling the argument of a Dirichlet series with a natural number, i. e. turning the series

$$A(s) = \sum_{n=1}^{\infty} \frac{a_n}{n^s}$$

into the series

$$A(ks) = \sum_{n=1}^{\infty} \frac{a_n}{(n^k)^2}.$$

definition *fds-scale* :: $\text{nat} \Rightarrow ('a :: \text{zero}) \text{fds} \Rightarrow 'a \text{fds}$ **where**
 $\text{fds-scale } c f =$
 $\text{fds } (\lambda n. \text{if } n > 0 \wedge \text{is-nth-power } c n \text{ then } \text{fds-nth } f (\text{nth-root-nat } c n) \text{ else } 0)$

lemma *fds-scale-0* [*simp*]: $\text{fds-scale } 0 f = 0$
<proof>

lemma *fds-scale-1* [*simp*]: $\text{fds-scale } 1 f = f$
<proof>

lemma *fds-nth-scale-power* [*simp*]:
 $c > 0 \implies \text{fds-nth } (\text{fds-scale } c f) (n \wedge c) = \text{fds-nth } f n$

<proof>

lemma *fds-nth-scale-nonpower* [*simp*]:

$\neg \text{is-nth-power } c \ n \implies \text{fds-nth } (\text{fds-scale } c \ f) \ n = 0$

<proof>

lemma *fds-nth-scale*:

$\text{fds-nth } (\text{fds-scale } c \ f) \ n =$

$(\text{if } n > 0 \wedge \text{is-nth-power } c \ n \text{ then } \text{fds-nth } f \ (\text{nth-root-nat } c \ n) \text{ else } 0)$

<proof>

lemma *fds-scale-const* [*simp*]: $c > 0 \implies \text{fds-scale } c \ (\text{fds-const } c') = \text{fds-const } c'$

<proof>

lemma *fds-scale-zero* [*simp*]: $\text{fds-scale } c \ 0 = 0$

<proof>

lemma *fds-scale-one* [*simp*]: $c > 0 \implies \text{fds-scale } c \ 1 = 1$

<proof>

lemma *fds-scale-of-nat* [*simp*]: $c > 0 \implies \text{fds-scale } c \ (\text{of-nat } n) = \text{of-nat } n$

<proof>

lemma *fds-scale-of-int* [*simp*]: $c > 0 \implies \text{fds-scale } c \ (\text{of-int } n) = \text{of-int } n$

<proof>

lemma *fds-scale-numeral* [*simp*]: $c > 0 \implies \text{fds-scale } c \ (\text{numeral } n) = \text{numeral } n$

<proof>

lemma *fds-scale-scale*: $\text{fds-scale } c \ (\text{fds-scale } c' \ f) = \text{fds-scale } (c * c') \ f$

<proof>

lemma *fds-scale-add* [*simp*]:

fixes $f \ g :: 'a :: \text{monoid-add } \text{fds}$

shows $\text{fds-scale } c \ (f + g) = \text{fds-scale } c \ f + \text{fds-scale } c \ g$

<proof>

lemma *fds-scale-minus* [*simp*]:

fixes $f \ g :: 'a :: \{\text{cancel-comm-monoid-add}\} \text{fds}$

shows $\text{fds-scale } c \ (f - g) = \text{fds-scale } c \ f - \text{fds-scale } c \ g$

<proof>

lemma *fds-scale-uminus* [*simp*]:

fixes $f :: 'a :: \text{group-add } \text{fds}$

shows $\text{fds-scale } c \ (-f) = -\text{fds-scale } c \ f$

<proof>

lemma *fds-scale-mult* [*simp*]:

fixes $f \ g :: 'a :: \text{semiring-0 } \text{fds}$

shows $\text{fds-scale } c (f * g) = \text{fds-scale } c f * \text{fds-scale } c g$
 ⟨proof⟩

lemma *fds-scale-shift*:

$\text{fds-shift } d (\text{fds-scale } c f) = \text{fds-scale } c (\text{fds-shift } (c * d) f)$
 ⟨proof⟩

lemma *fds-ind-nth-power*: $k > 0 \implies \text{fds-ind } (\text{is-nth-power } k) = \text{fds-scale } k \text{ fds-zeta}$
 ⟨proof⟩

4.4 Formal derivative

The formal derivative of a series

$$A(s) = \sum_{n=1}^{\infty} \frac{a_n}{n^s}$$

can easily be seen to be

$$A'(s) = - \sum_{n=1}^{\infty} \frac{\ln n \cdot a_n}{n^s} .$$

definition *fds-deriv* :: 'a :: real-algebra $\text{fds} \Rightarrow$ 'a *fds* **where**
 $\text{fds-deriv } f = \text{fds } (\lambda n. - \ln (\text{real } n) *_{\mathbb{R}} \text{fds-nth } f n)$

lemma *fds-nth-deriv*: $\text{fds-nth } (\text{fds-deriv } f) n = -\ln (\text{real } n) *_{\mathbb{R}} \text{fds-nth } f n$
 ⟨proof⟩

lemma *fds-deriv-const* [simp]: $\text{fds-deriv } (\text{fds-const } c) = 0$
 ⟨proof⟩

lemma *fds-deriv-0* [simp]: $\text{fds-deriv } 0 = 0$
 ⟨proof⟩

lemma *fds-deriv-1* [simp]: $\text{fds-deriv } 1 = 0$
 ⟨proof⟩

lemma *fds-deriv-of-nat* [simp]: $\text{fds-deriv } (\text{of-nat } n) = 0$
 ⟨proof⟩

lemma *fds-deriv-of-int* [simp]: $\text{fds-deriv } (\text{of-int } n) = 0$
 ⟨proof⟩

lemma *fds-deriv-of-real* [simp]: $\text{fds-deriv } (\text{of-real } n) = 0$
 ⟨proof⟩

lemma *fds-deriv-uminus* [simp]: $\text{fds-deriv } (-f) = -\text{fds-deriv } f$
 ⟨proof⟩

lemma *fds-deriv-add* [*simp*]: $fds\text{-deriv } (f + g) = fds\text{-deriv } f + fds\text{-deriv } g$
 ⟨*proof*⟩

lemma *fds-deriv-minus* [*simp*]: $fds\text{-deriv } (f - g) = fds\text{-deriv } f - fds\text{-deriv } g$
 ⟨*proof*⟩

lemma *fds-deriv-times* [*simp*]:
 $fds\text{-deriv } (f * g) = fds\text{-deriv } f * g + f * fds\text{-deriv } g$
 ⟨*proof*⟩

lemma *fds-deriv-inverse* [*simp*]:
fixes $f :: 'a :: \{real\text{-algebra}, field\}$ *fds*
assumes $fds\text{-nth } f (Suc\ 0) \neq 0$
shows $fds\text{-deriv } (inverse\ f) = -fds\text{-deriv } f / f^{\wedge} 2$
 ⟨*proof*⟩

lemma *fds-deriv-shift* [*simp*]: $fds\text{-deriv } (fds\text{-shift } c\ f) = fds\text{-shift } c\ (fds\text{-deriv } f)$
 ⟨*proof*⟩

lemma *fds-deriv-scale*: $fds\text{-deriv } (fds\text{-scale } c\ f) = of\text{-nat } c * fds\text{-scale } c\ (fds\text{-deriv } f)$
 ⟨*proof*⟩

lemma *fds-deriv-eq-imp-eq*:
assumes $fds\text{-deriv } f = fds\text{-deriv } g$ $fds\text{-nth } f (Suc\ 0) = fds\text{-nth } g (Suc\ 0)$
shows $f = g$
 ⟨*proof*⟩

lemma *completely-multiplicative-fds-deriv*:
assumes *completely-multiplicative-function* f
shows $fds\text{-deriv } (fds\ f) = -fds\ (\lambda n. f\ n * mangoldt\ n) * fds\ f$
 ⟨*proof*⟩

lemma *completely-multiplicative-fds-deriv'*:
completely-multiplicative-function $(fds\text{-nth } f) \implies$
 $fds\text{-deriv } f = -fds\ (\lambda n. fds\text{-nth } f\ n * mangoldt\ n) * f$
 ⟨*proof*⟩

lemma *fds-deriv-zeta*:
 $fds\text{-deriv } fds\text{-zeta} =$
 $-fds\ mangoldt * (fds\text{-zeta} :: 'a :: \{comm\text{-semiring-1}, real\text{-algebra-1}\})\ fds$
 ⟨*proof*⟩

lemma *fds-mangoldt-times-zeta*: $fds\ mangoldt * fds\text{-zeta} = fds\ (\lambda x. of\text{-real } (ln\ (real\ x)))$
 ⟨*proof*⟩

lemma *fds-deriv-zeta'*: $fds\text{-deriv } fds\text{-zeta} =$

$-fds$ ($\lambda x. of\text{-}real$ (ln ($real$ x))) :: 'a :: {*comm-semiring-1, real-algebra-1*}

<proof>

4.5 Formal integral

definition *fds-integral* :: 'a \Rightarrow 'a :: *real-algebra* *fds* \Rightarrow 'a *fds* **where**
fds-integral c $f = fds$ ($\lambda n. if$ $n = 1$ *then* c *else* $- fds\text{-}nth$ f $n /_R ln$ (*real* n))

lemma *fds-integral-0* [*simp*]: *fds-integral* a $0 = fds\text{-}const$ a
<proof>

lemma *fds-integral-add*: *fds-integral* ($a + b$) ($f + g$) = *fds-integral* a $f + fds\text{-}integral$ b g
<proof>

lemma *fds-integral-diff*: *fds-integral* ($a - b$) ($f - g$) = *fds-integral* a $f - fds\text{-}integral$ b g
<proof>

lemma *fds-integral-minus*: *fds-integral* ($-a$) ($-f$) = $-fds\text{-}integral$ a f
<proof>

lemma *fds-shift-integral*: *fds-shift* b (*fds-integral* a f) = *fds-integral* a (*fds-shift* b f)
<proof>

lemma *fds-deriv-fds-integral* [*simp*]:
fds-nth f (*Suc* 0) = $0 \implies fds\text{-}deriv$ (*fds-integral* c f) = f
<proof>

lemma *fds-integral-fds-deriv* [*simp*]: *fds-integral* (*fds-nth* f 1) (*fds-deriv* f) = f
<proof>

4.6 Formal logarithm

definition *fds-ln* :: 'a \Rightarrow 'a :: {*real-normed-field*} *fds* \Rightarrow 'a *fds* **where**
fds-ln l $f = fds\text{-}integral$ l (*fds-deriv* f / f)

lemma *fds-nth-Suc-0-fds-deriv* [*simp*]: *fds-nth* (*fds-deriv* f) (*Suc* 0) = 0
<proof>

lemma *fds-deriv-fds-ln* [*simp*]: *fds-deriv* (*fds-ln* l f) = *fds-deriv* f / f
<proof>

lemma *fds-nth-Suc-0-fds-ln* [*simp*]: *fds-nth* (*fds-ln* l f) (*Suc* 0) = l
<proof>

lemma *fds-ln-const* [*simp*]: *fds-ln* l (*fds-const* c) = *fds-const* l
<proof>

lemma *fds-ln-0* [*simp*]: $\text{fds-ln } l \ 0 = \text{fds-const } l$
 ⟨*proof*⟩

lemma *fds-ln-1* [*simp*]: $\text{fds-ln } l \ 1 = \text{fds-const } l$
 ⟨*proof*⟩

lemma *fds-shift-ln* [*simp*]: $\text{fds-shift } a \ (\text{fds-ln } l \ f) = \text{fds-ln } l \ (\text{fds-shift } a \ f)$
 ⟨*proof*⟩

lemma *fds-ln-mult*:
 assumes $\text{fds-nth } f \ 1 \neq 0 \ \text{fds-nth } g \ 1 \neq 0 \ l' + l'' = l$
 shows $\text{fds-ln } l \ (f * g) = \text{fds-ln } l' \ f + \text{fds-ln } l'' \ g$
 ⟨*proof*⟩

lemma *fds-ln-power*:
 assumes $\text{fds-nth } f \ 1 \neq 0 \ l = \text{of-nat } n * l'$
 shows $\text{fds-ln } l \ (f \wedge n) = \text{of-nat } n * \text{fds-ln } l' \ f$
 ⟨*proof*⟩

lemma *fds-ln-prod*:
 assumes $\bigwedge x. x \in A \implies \text{fds-nth } (f \ x) \ 1 \neq 0 \ (\sum x \in A. l' \ x) = l$
 shows $\text{fds-ln } l \ (\prod x \in A. f \ x) = (\sum x \in A. \text{fds-ln } (l' \ x) \ (f \ x))$
 ⟨*proof*⟩

4.7 Formal exponential

definition *fds-exp* :: 'a :: {*real-normed-algebra-1,banach*} *fds* \Rightarrow 'a *fds* **where**
 $\text{fds-exp } f = (\text{let } f' = \text{fds } (\lambda n. \text{if } n = 1 \text{ then } 0 \text{ else } \text{fds-nth } f \ n)$
 $\text{in } \text{fds } (\lambda n. \text{exp } (\text{fds-nth } f \ 1) * (\sum k. \text{fds-nth } (f' \wedge k) \ n \ /_R \ \text{fact } k)))$

lemma *fds-nth-exp-Suc-0* [*simp*]: $\text{fds-nth } (\text{fds-exp } f) \ (\text{Suc } 0) = \text{exp } (\text{fds-nth } f \ 1)$
 ⟨*proof*⟩

lemma *fds-exp-times-fds-nth-0*:
 $\text{fds-const } (\text{exp } (\text{fds-nth } f \ (\text{Suc } 0))) * \text{fds-exp } (f - \text{fds-const } (\text{fds-nth } f \ (\text{Suc } 0)))$
 $= \text{fds-exp } f$
 ⟨*proof*⟩

lemma *fds-exp-const* [*simp*]: $\text{fds-exp } (\text{fds-const } c) = \text{fds-const } (\text{exp } c)$
 ⟨*proof*⟩

lemma *fds-exp-numeral* [*simp*]: $\text{fds-exp } (\text{numeral } n) = \text{fds-const } (\text{exp } (\text{numeral } n))$
 ⟨*proof*⟩

lemma *fds-exp-0* [*simp*]: $\text{fds-exp } 0 = 1$
 ⟨*proof*⟩

lemma *fds-exp-1* [*simp*]: $\text{fds-exp } 1 = \text{fds-const } (\text{exp } 1)$
 ⟨*proof*⟩

lemma *fds-nth-Suc-0-exp* [simp]: $fds_nth (fds_exp f) (Suc 0) = exp (fds_nth f (Suc 0))$
 ⟨proof⟩

4.8 Subseries

definition *fds-subseries* :: $(nat \Rightarrow bool) \Rightarrow ('a :: semiring-1) fds \Rightarrow 'a fds$ **where**
 $fds_subseries P f = fds (\lambda n. if P n then fds_nth f n else 0)$

lemma *fds-nth-subseries*:
 $fds_nth (fds_subseries P f) n = (if P n then fds_nth f n else 0)$
 ⟨proof⟩

lemma *fds-subseries-0* [simp]: $fds_subseries P 0 = 0$
 ⟨proof⟩

lemma *fds-subseries-1* [simp]: $P 1 \Longrightarrow fds_subseries P 1 = 1$
 ⟨proof⟩

lemma *fds-subseries-const* [simp]: $P 1 \Longrightarrow fds_subseries P (fds_const c) = fds_const c$
 ⟨proof⟩

lemma *fds-subseries-add* [simp]: $fds_subseries P (f + g) = fds_subseries P f + fds_subseries P g$
 ⟨proof⟩

lemma *fds-subseries-diff* [simp]:
 $fds_subseries P (f - g :: 'a :: ring-1 fds) = fds_subseries P f - fds_subseries P g$
 ⟨proof⟩

lemma *fds-subseries-minus* [simp]:
 $fds_subseries P (-f :: 'a :: ring-1 fds) = - fds_subseries P f$
 ⟨proof⟩

lemma *fds-subseries-sum* [simp]: $fds_subseries P (\sum x \in A. f x) = (\sum x \in A. fds_subseries P (f x))$
 ⟨proof⟩

lemma *fds-subseries-shift* [simp]:
 $fds_subseries P (fds_shift c f) = fds_shift c (fds_subseries P f)$
 ⟨proof⟩

lemma *fds-subseries-deriv* [simp]:
 $fds_subseries P (fds_deriv f) = fds_deriv (fds_subseries P f)$
 ⟨proof⟩

lemma *fds-subseries-integral* [simp]:

$P \ 1 \vee c = 0 \implies \text{fds-subseries } P \ (\text{fds-integral } c \ f) = \text{fds-integral } c \ (\text{fds-subseries } P \ f)$
 ⟨proof⟩

abbreviation $\text{fds-primelow-subseries} :: \text{nat} \Rightarrow ('a :: \text{semiring-1}) \text{fds} \Rightarrow 'a \ \text{fds}$
where

$\text{fds-primelow-subseries } p \ f \equiv \text{fds-subseries } (\lambda n. \text{prime-factors } n \subseteq \{p\}) \ f$

lemma $\text{fds-primelow-subseries-mult}$ [simp]:

fixes $p :: \text{nat}$

defines $P \equiv (\lambda n. \text{prime-factors } n \subseteq \{p\})$

shows $\text{fds-subseries } P \ (f * g) = \text{fds-subseries } P \ f * \text{fds-subseries } P \ g$

⟨proof⟩

lemma $\text{fds-primelow-subseries-power}$ [simp]:

$\text{fds-primelow-subseries } p \ (f \wedge n) = \text{fds-primelow-subseries } p \ f \wedge n$

⟨proof⟩

lemma $\text{fds-primelow-subseries-prod}$ [simp]:

$\text{fds-primelow-subseries } p \ (\prod x \in A. f \ x) = (\prod x \in A. \text{fds-primelow-subseries } p \ (f \ x))$

⟨proof⟩

lemma $\text{completely-multiplicative-function-only-pows}$:

assumes $\text{completely-multiplicative-function } (\text{fds-nth } f)$

shows $\text{completely-multiplicative-function } (\text{fds-nth } (\text{fds-primelow-subseries } p \ f))$

⟨proof⟩

4.9 Truncation

definition $\text{fds-truncate} :: \text{nat} \Rightarrow 'a :: \{\text{zero}\} \ \text{fds} \Rightarrow 'a \ \text{fds}$ **where**

$\text{fds-truncate } m \ f = \text{fds } (\lambda n. \text{if } n \leq m \text{ then } \text{fds-nth } f \ n \ \text{else } 0)$

lemma fds-nth-truncate : $\text{fds-nth } (\text{fds-truncate } m \ f) \ n = (\text{if } n \leq m \text{ then } \text{fds-nth } f \ n \ \text{else } 0)$

⟨proof⟩

lemma fds-truncate-0 [simp]: $\text{fds-truncate } 0 \ f = 0$

⟨proof⟩

lemma fds-truncate-zero [simp]: $\text{fds-truncate } m \ 0 = 0$

⟨proof⟩

lemma fds-truncate-one [simp]: $m > 0 \implies \text{fds-truncate } m \ 1 = 1$

⟨proof⟩

lemma $\text{fds-truncate-const}$ [simp]: $m > 0 \implies \text{fds-truncate } m \ (\text{fds-const } c) = \text{fds-const } c$

⟨proof⟩

lemma *fds-truncate-truncate* [simp]: $fds-truncate\ m\ (fds-truncate\ n\ f) = fds-truncate\ (min\ m\ n)\ f$
 ⟨proof⟩

lemma *fds-truncate-truncate'* [simp]: $fds-truncate\ m\ (fds-truncate\ m\ f) = fds-truncate\ m\ f$
 ⟨proof⟩

lemma *fds-truncate-shift* [simp]: $fds-truncate\ m\ (fds-shift\ a\ f) = fds-shift\ a\ (fds-truncate\ m\ f)$
 ⟨proof⟩

lemma *fds-truncate-add-strong*:
 $fds-truncate\ m\ (f + g :: 'a :: monoid-add\ fds) = fds-truncate\ m\ f + fds-truncate\ m\ g$
 ⟨proof⟩

lemma *fds-truncate-add*:
 $fds-truncate\ m\ (fds-truncate\ m\ f + fds-truncate\ m\ g :: 'a :: monoid-add\ fds) =$
 $fds-truncate\ m\ (f + g)$
 ⟨proof⟩

lemma *fds-truncate-mult*:
 $fds-truncate\ m\ (fds-truncate\ m\ f * fds-truncate\ m\ g) = fds-truncate\ m\ (f * g)$ (is
 ?A = ?B)
 ⟨proof⟩

lemma *fds-truncate-deriv*: $fds-truncate\ m\ (fds-deriv\ f) = fds-deriv\ (fds-truncate\ m\ f)$
 ⟨proof⟩

lemma *fds-truncate-integral*:
 $m > 0 \vee c = 0 \implies fds-truncate\ m\ (fds-integral\ c\ f) = fds-integral\ c\ (fds-truncate\ m\ f)$
 ⟨proof⟩

lemma *fds-truncate-power*: $fds-truncate\ m\ (fds-truncate\ m\ f ^ n) = fds-truncate\ m\ (f ^ n)$
 ⟨proof⟩

lemma *dirichlet-inverse-cong-simp*:
assumes $\bigwedge m. m > 0 \implies m \leq n \implies f\ m = f'\ m\ i = i'\ n = n'$
shows $dirichlet-inverse\ f\ i\ n = dirichlet-inverse\ f'\ i'\ n'$
 ⟨proof⟩

lemma *fds-truncate-cong*:
 $(\bigwedge n. m > 0 \implies n > 0 \implies n \leq m \implies fds-nth\ f\ n = fds-nth\ f'\ n) \implies$
 $fds-truncate\ m\ f = fds-truncate\ m\ f'$

<proof>

lemma *fds-truncate-inverse:*

fds-truncate m (inverse (fds-truncate m (f :: 'a :: field fds))) = fds-truncate m (inverse f)
<proof>

lemma *fds-truncate-divide:*

fixes *f g :: 'a :: field fds*
shows *fds-truncate m (fds-truncate m f / fds-truncate m g) = fds-truncate m (f / g)*
<proof>

lemma *fds-truncate-ln:*

fixes *f :: 'a :: real-normed-field fds*
shows *fds-truncate m (fds-ln l (fds-truncate m f)) = fds-truncate m (fds-ln l f)*
<proof>

lemma *fds-truncate-exp:*

shows *fds-truncate m (fds-exp (fds-truncate m f)) = fds-truncate m (fds-exp f)*
<proof>

lemma *fds-eqI-truncate:*

assumes $\bigwedge m. m > 0 \implies \text{fds-truncate } m \ f = \text{fds-truncate } m \ g$
shows $f = g$
<proof>

4.10 Normed series

definition *fds-norm :: 'a :: {real-normed-div-algebra} fds \Rightarrow real fds*
where *fds-norm f = fds ($\lambda n. \text{of-real (norm (fds-nth f n))}$)*

lemma *fds-nth-norm [simp]: fds-nth (fds-norm f) n = norm (fds-nth f n)*
<proof>

lemma *fds-norm-1 [simp]: fds-norm 1 = 1*
<proof>

lemma *fds-nth-norm-mult-le:*

shows $\text{norm (fds-nth (f * g) n)} \leq \text{fds-nth (fds-norm f * fds-norm g) n}$
<proof>

lemma *fds-nth-norm-mult-nonneg [simp]: fds-nth (fds-norm f * fds-norm g) n \geq 0*
<proof>

4.11 Lifting a real series to a real algebra

definition *fds-of-real :: real fds \Rightarrow 'a :: {real-normed-algebra-1} fds* **where**
fds-of-real f = fds ($\lambda n. \text{of-real (fds-nth f n)}$)

lemma *fds-nth-of-real* [*simp*]: $fds\text{-}nth\ (fds\text{-}of\text{-}real\ f)\ n = of\text{-}real\ (fds\text{-}nth\ f\ n)$
 ⟨*proof*⟩

lemma *fds-of-real-0* [*simp*]: $fds\text{-}of\text{-}real\ 0 = 0$
and *fds-of-real-1* [*simp*]: $fds\text{-}of\text{-}real\ 1 = 1$
and *fds-of-real-const* [*simp*]: $fds\text{-}of\text{-}real\ (fds\text{-}const\ c) = fds\text{-}const\ (of\text{-}real\ c)$
and *fds-of-real-minus* [*simp*]: $fds\text{-}of\text{-}real\ (-f) = -fds\text{-}of\text{-}real\ f$
and *fds-of-real-add* [*simp*]: $fds\text{-}of\text{-}real\ (f + g) = fds\text{-}of\text{-}real\ f + fds\text{-}of\text{-}real\ g$
and *fds-of-real-mult* [*simp*]: $fds\text{-}of\text{-}real\ (f * g) = fds\text{-}of\text{-}real\ f * fds\text{-}of\text{-}real\ g$
and *fds-of-real-deriv* [*simp*]: $fds\text{-}of\text{-}real\ (fds\text{-}deriv\ f) = fds\text{-}deriv\ (fds\text{-}of\text{-}real\ f)$
 ⟨*proof*⟩

lemma *fds-of-real-higher-deriv* [*simp*]:
 $(fds\text{-}deriv\ \widehat{\widehat{n}}\ (fds\text{-}of\text{-}real\ f)) = fds\text{-}of\text{-}real\ ((fds\text{-}deriv\ \widehat{\widehat{n}})\ f)$
 ⟨*proof*⟩

4.12 Convergence and connection to concrete functions

The following definitions establish a connection of a formal Dirichlet series to the concrete analytic function that it corresponds to. This correspondence is usually partial in the sense that a series may not converge everywhere.

definition *eval-fds* :: $('a :: \{nat\text{-}power, real\text{-}normed\text{-}field, banach\})\ fds \Rightarrow 'a \Rightarrow 'a$
where

$$eval\text{-}fds\ f\ s = (\sum n. fds\text{-}nth\ f\ n / nat\text{-}power\ n\ s)$$

lemma *eval-fds-eqI*:

assumes $(\lambda n. fds\text{-}nth\ f\ (Suc\ n) / nat\text{-}power\ (Suc\ n)\ s)\ sums\ L$

shows $eval\text{-}fds\ f\ s = L$

⟨*proof*⟩

definition *fds-converges* ::

$('a :: \{nat\text{-}power, real\text{-}normed\text{-}field, banach\})\ fds \Rightarrow 'a \Rightarrow bool$ **where**
 $fds\text{-}converges\ f\ s \iff summable\ (\lambda n. fds\text{-}nth\ f\ n / nat\text{-}power\ n\ s)$

lemma *fds-converges-iff*:

$fds\text{-}converges\ f\ s \iff (\lambda n. fds\text{-}nth\ f\ n / nat\text{-}power\ n\ s)\ sums\ eval\text{-}fds\ f\ s$

⟨*proof*⟩

definition *fds-abs-converges* ::

$('a :: \{nat\text{-}power, real\text{-}normed\text{-}field, banach\})\ fds \Rightarrow 'a \Rightarrow bool$ **where**
 $fds\text{-}abs\text{-}converges\ f\ s \iff summable\ (\lambda n. norm\ (fds\text{-}nth\ f\ n / nat\text{-}power\ n\ s))$

lemma *fds-abs-converges-imp-converges* [*dest*, *intro*]:

$fds\text{-}abs\text{-}converges\ f\ s \implies fds\text{-}converges\ f\ s$

⟨*proof*⟩

lemma *fds-converges-altdef*:

fds-converges $f\ s \iff (\lambda n. \text{fds-nth } f \ (\text{Suc } n) / \text{nat-power } (\text{Suc } n) \ s) \text{ sums eval-fds } f\ s$
 ⟨proof⟩

lemma *fds-const-abs-converges* [simp]: *fds-abs-converges* (*fds-const* c) s
 ⟨proof⟩

lemma *fds-const-converges* [simp]: *fds-converges* (*fds-const* c) s
 ⟨proof⟩

lemma *eval-fds-const* [simp]: *eval-fds* (*fds-const* c) = $(\lambda-. c)$
 ⟨proof⟩

lemma *fds-zero-abs-converges* [simp]: *fds-abs-converges* $0\ s$
 ⟨proof⟩

lemma *fds-zero-converges* [simp]: *fds-converges* $0\ s$
 ⟨proof⟩

lemma *eval-fds-zero* [simp]: *eval-fds* $0 = (\lambda-. 0)$
 ⟨proof⟩

lemma *fds-one-abs-converges* [simp]: *fds-abs-converges* $1\ s$
 ⟨proof⟩

lemma *fds-one-converges* [simp]: *fds-converges* $1\ s$
 ⟨proof⟩

lemma *fds-converges-truncate* [simp]: *fds-converges* (*fds-truncate* $n\ f$) s
 ⟨proof⟩

lemma *fds-abs-converges-truncate* [simp]: *fds-abs-converges* (*fds-truncate* $n\ f$) s
 ⟨proof⟩

lemma *fds-abs-converges-subseries* [simp, intro]:
assumes *fds-abs-converges* $f\ s$
shows *fds-abs-converges* (*fds-subseries* $P\ f$) s
 ⟨proof⟩

lemma *eval-fds-one* [simp]: *eval-fds* $1 = (\lambda-. 1)$
 ⟨proof⟩

lemma *eval-fds-truncate*: *eval-fds* (*fds-truncate* $n\ f$) $s = (\sum k=1..n. \text{fds-nth } f\ k / \text{nat-power } k\ s)$
 ⟨proof⟩

lemma *fds-converges-add*:
assumes *fds-converges* $f\ s$ *fds-converges* $g\ s$

shows $fds\text{-converges } (f + g) s$
 $\langle proof \rangle$

lemma $fds\text{-abs-converges-add}$:
assumes $fds\text{-abs-converges } f s$ $fds\text{-abs-converges } g s$
shows $fds\text{-abs-converges } (f + g) s$
 $\langle proof \rangle$

lemma $eval\text{-fds-add}$:
assumes $fds\text{-converges } f s$ $fds\text{-converges } g s$
shows $eval\text{-fds } (f + g) s = eval\text{-fds } f s + eval\text{-fds } g s$
 $\langle proof \rangle$

lemma $fds\text{-converges-uminus}$:
assumes $fds\text{-converges } f s$
shows $fds\text{-converges } (-f) s$
 $\langle proof \rangle$

lemma $The\text{-cong}$: $The P = The Q$ if $\bigwedge x. P x \longleftrightarrow Q x$
 $\langle proof \rangle$

lemma $fds\text{-abs-converges-uminus}$:
assumes $fds\text{-abs-converges } f s$
shows $fds\text{-abs-converges } (-f) s$
 $\langle proof \rangle$

lemma $eval\text{-fds-uminus}$: $fds\text{-converges } f s \implies eval\text{-fds } (-f) s = -eval\text{-fds } f s$
 $\langle proof \rangle$

lemma $fds\text{-converges-diff}$:
assumes $fds\text{-converges } f s$ $fds\text{-converges } g s$
shows $fds\text{-converges } (f - g) s$
 $\langle proof \rangle$

lemma $fds\text{-abs-converges-diff}$:
assumes $fds\text{-abs-converges } f s$ $fds\text{-abs-converges } g s$
shows $fds\text{-abs-converges } (f - g) s$
 $\langle proof \rangle$

lemma $eval\text{-fds-diff}$:
assumes $fds\text{-converges } f s$ $fds\text{-converges } g s$
shows $eval\text{-fds } (f - g) s = eval\text{-fds } f s - eval\text{-fds } g s$
 $\langle proof \rangle$

lemma $eval\text{-fds-at-nat}$: $eval\text{-fds } f (of\text{-nat } k) = (\sum n. fds\text{-nth } f n / of\text{-nat } n \wedge k)$
 $\langle proof \rangle$

lemma *eval-fds-at-numeral*: $\text{eval-fds } f \text{ (numeral } k) = (\sum n. \text{fds-nth } f \text{ } n / \text{of-nat } n \hat{\text{ numeral } k})$
 ⟨proof⟩

lemma *eval-fds-at-1*: $\text{eval-fds } f \text{ } 1 = (\sum n. \text{fds-nth } f \text{ } n / \text{of-nat } n)$
 ⟨proof⟩

lemma *eval-fds-at-0*: $\text{eval-fds } f \text{ } 0 = (\sum n. \text{fds-nth } f \text{ } n)$
 ⟨proof⟩

lemma *suminf-fds-zeta-aux*:
 $f \text{ } 0 = 0 \implies (\sum n. \text{fds-nth } \text{fds-zeta } n / f \text{ } n) = (\sum n. 1 / f \text{ } n :: 'a :: \text{real-normed-field})$
 ⟨proof⟩

lemma *fds-converges-shift [simp]*:
fixes $z :: 'a :: \{\text{banach, nat-power-field, real-normed-field}\}$
shows $\text{fds-converges } (\text{fds-shift } c \text{ } f) \text{ } z \longleftrightarrow \text{fds-converges } f \text{ } (z - c)$
 ⟨proof⟩

lemma *fds-abs-converges-shift [simp]*:
fixes $z :: 'a :: \{\text{banach, nat-power-field, real-normed-field}\}$
shows $\text{fds-abs-converges } (\text{fds-shift } c \text{ } f) \text{ } z \longleftrightarrow \text{fds-abs-converges } f \text{ } (z - c)$
 ⟨proof⟩

lemma *fds-eval-shift [simp]*:
fixes $z :: 'a :: \{\text{banach, nat-power-field, real-normed-field}\}$
shows $\text{eval-fds } (\text{fds-shift } c \text{ } f) \text{ } z = \text{eval-fds } f \text{ } (z - c)$
 ⟨proof⟩

lemma *fds-converges-scale [simp]*:
fixes $z :: 'a :: \{\text{banach, nat-power-field, real-normed-field}\}$
assumes $c: c > 0$
shows $\text{fds-converges } (\text{fds-scale } c \text{ } f) \text{ } z \longleftrightarrow \text{fds-converges } f \text{ } (\text{of-nat } c * z)$
 ⟨proof⟩

lemma *fds-abs-converges-scale [simp]*:
fixes $z :: 'a :: \{\text{banach, nat-power-field, real-normed-field}\}$
assumes $c: c > 0$
shows $\text{fds-abs-converges } (\text{fds-scale } c \text{ } f) \text{ } z \longleftrightarrow \text{fds-abs-converges } f \text{ } (\text{of-nat } c * z)$
 ⟨proof⟩

lemma *eval-fds-scale [simp]*:
fixes $z :: 'a :: \{\text{banach, nat-power-field, real-normed-field}\}$
assumes $c: c > 0$
shows $\text{eval-fds } (\text{fds-scale } c \text{ } f) \text{ } z = \text{eval-fds } f \text{ } (\text{of-nat } c * z)$
 ⟨proof⟩

lemma *fds-abs-converges-integral*:
assumes *fds-abs-converges f s*
shows *fds-abs-converges (fds-integral c f) s*
<proof>

lemma *fds-abs-converges-ln*:
assumes *fds-abs-converges (fds-deriv f / f) s*
shows *fds-abs-converges (fds-ln l f) s*
<proof>

end

5 The Möbius μ function

theory *Moebius-Mu*

imports

Main

HOL-Number-Theory.Number-Theory

HOL-Computational-Algebra.Squarefree

Dirichlet-Series

Dirichlet-Misc

begin

definition *moebius-mu* :: *nat* \Rightarrow *'a* :: *comm-ring-1* **where**
moebius-mu n =
(if squarefree n then (-1) ^ card (prime-factors n) else 0)

lemma *abs-moebius-mu-le*: *abs (moebius-mu n :: 'a :: {linordered-idom})* ≤ 1
<proof>

lemma *of-int-moebius-mu [simp]*: *of-int (moebius-mu n)* = *moebius-mu n*
<proof>

lemma *minus-1-power-ring-neq-zero [simp]*: *(- 1 :: 'a :: ring-1)* $^ n \neq 0$
<proof>

lemma *moebius-mu-0 [simp]*: *moebius-mu 0* = *0*
<proof>

lemma *fds-nth-fds-moebius-mu [simp]*: *fds-nth (fds moebius-mu)* = *moebius-mu*
<proof>

lemma *prime-factors-Suc-0 [simp]*: *prime-factors (Suc 0)* = *{}*
<proof>

lemma *moebius-mu-Suc-0 [simp]*: *moebius-mu (Suc 0)* = *1*
<proof>

lemma *moebius-mu-1* [*simp*]: $\text{moebius-mu } 1 = 1$
{*proof*}

lemma *moebius-mu-eq-zero-iff*: $\text{moebius-mu } n = 0 \iff \neg \text{squarefree } n$
{*proof*}

lemma *moebius-mu-not-squarefree* [*simp*]: $\neg \text{squarefree } n \implies \text{moebius-mu } n = 0$
{*proof*}

lemma *moebius-mu-power*:
 assumes $a > 1 \ n > 1$
 shows $\text{moebius-mu } (a \wedge n) = 0$
{*proof*}

lemma *moebius-mu-power'*:
 $\text{moebius-mu } (a \wedge n) = (\text{if } a = 1 \vee n = 0 \text{ then } 1 \text{ else if } n = 1 \text{ then } \text{moebius-mu } a \text{ else } 0)$
{*proof*}

lemma *moebius-mu-squarefree-eq*:
 $\text{squarefree } n \implies \text{moebius-mu } n = (-1) \wedge \text{card } (\text{prime-factors } n)$
{*proof*}

lemma *moebius-mu-squarefree-eq'*:
 assumes $\text{squarefree } n$
 shows $\text{moebius-mu } n = (-1) \wedge \text{size } (\text{prime-factorization } n)$
{*proof*}

lemma *sum-moebius-mu-divisors*:
 assumes $n > 1$
 shows $(\sum d \mid d \text{ dvd } n. \text{moebius-mu } d) = (0 \text{ :: 'a :: comm-ring-1})$
{*proof*}

lemma *sum-moebius-mu-divisors'*:
 $(\sum d \mid d \text{ dvd } n. \text{moebius-mu } d) = (\text{if } n = 1 \text{ then } 1 \text{ else } 0)$
{*proof*}

lemma *fds-zeta-times-moebius-mu*: $\text{fds-zeta} * \text{fds moebius-mu} = 1$
{*proof*}

lemma *fds-moebius-inverse-zeta*:
 $\text{fds moebius-mu} = \text{inverse } (\text{fds-zeta} \text{ :: 'a :: field } \text{fds})$
{*proof*}

lemma *moebius-mu-formula-real*: $(\text{moebius-mu } n \text{ :: real}) = \text{dirichlet-inverse } (\lambda-. 1) \ 1 \ n$
{*proof*}

lemma *moebius-mu-formula-int*: $\text{moebius-mu } n = \text{dirichlet-inverse } (\lambda-. 1 \text{ :: int}) \ 1$

n
<proof>

lemma *moebius-mu-formula*: $\text{moebius-mu } n = \text{dirichlet-inverse } (\lambda-. 1) 1 n$
<proof>

interpretation *moebius-mu*: *multiplicative-function moebius-mu*
<proof>

interpretation *moebius-mu*:
multiplicative-function' moebius-mu $\lambda p k$. *if* $k = 1$ *then* -1 *else* 0 $\lambda-. -1$
<proof>

lemma *moebius-mu-2* [*simp*]: $\text{moebius-mu } 2 = -1$
and *moebius-mu-3* [*simp*]: $\text{moebius-mu } 3 = -1$
<proof>

lemma *moebius-mu-code* [*code*]:
 $\text{moebius-mu } n = \text{of-int } (\text{dirichlet-inverse } (\lambda-. 1 :: \text{int}) 1 n)$
<proof>

lemma *fds-moebius-inversion*: $f = \text{fds moebius-mu} * g \longleftrightarrow g = f * \text{fds-zeta}$
<proof>

lemma *moebius-inversion*:
assumes $\bigwedge n. n > 0 \implies g n = (\sum d \mid d \text{ dvd } n. f d) n > 0$
shows $f n = \text{dirichlet-prod moebius-mu } g n$
<proof>

lemma *fds-mangoldt*: $\text{fds mangoldt} = \text{fds moebius-mu} * \text{fds } (\lambda n. \text{of-real } (\ln (\text{real } n)))$
<proof>

lemma *sum-divisors-moebius-mu-times-multiplicative*:
fixes $f :: \text{nat} \Rightarrow 'a :: \{\text{comm-ring-1}\}$
assumes *multiplicative-function* $f n > 0$
shows $(\sum d \mid d \text{ dvd } n. \text{moebius-mu } d * f d) = (\prod p \in \text{prime-factors } n. 1 - f p)$
<proof>

lemma *completely-multiplicative-iff-inverse-moebius-mu*:
fixes $f :: \text{nat} \Rightarrow 'a :: \{\text{comm-ring-1}, \text{ring-no-zero-divisors}\}$
assumes *multiplicative-function* f
defines $g \equiv \text{dirichlet-inverse } f 1$
shows *completely-multiplicative-function* $f \longleftrightarrow$

$(\forall n. g\ n = \text{moebius-mu}\ n * f\ n)$
 <proof>

lemma *completely-multiplicative-fds-inverse*:
 fixes $f :: \text{nat} \Rightarrow 'a :: \text{field}$
 assumes *completely-multiplicative-function* f
 shows $\text{inverse}\ (f\text{ds}\ f) = f\text{ds}\ (\lambda n. \text{moebius-mu}\ n * f\ n)$
 <proof>

lemma *completely-multiplicative-fds-inverse'*:
 fixes $f :: 'a :: \text{field}\ \text{fds}$
 assumes *completely-multiplicative-function* $(f\text{ds-nth}\ f)$
 shows $\text{inverse}\ f = f\text{ds}\ (\lambda n. \text{moebius-mu}\ n * f\text{ds-nth}\ f\ n)$
 <proof>

context
 includes *fds-syntax*
 begin

lemma *selberg-aux*:
 $(\chi\ n. \text{of-real}\ ((\ln\ n)^2)) * f\text{ds}\ \text{moebius-mu} =$
 $(f\text{ds}\ \text{mangoldt})^2 - f\text{ds-deriv}\ (f\text{ds}\ \text{mangoldt} :: 'a :: \{\text{comm-ring-1}, \text{real-algebra-1}\})$
 $f\text{ds}$
 <proof>

lemma *selberg-aux'*:
 $\text{mangoldt}\ n * \text{of-real}\ (\ln\ n) + (\text{mangoldt} * \text{mangoldt})\ n =$
 $((\text{moebius-mu} * (\lambda b. \text{of-real}\ (\ln\ b)^2))\ n)$
 $:: 'a :: \{\text{comm-ring-1}, \text{real-algebra-1}\}$ if $n > 0$
 <proof>

end

end

6 Euler's ϕ function

theory *More-Totient*
 imports
 Moebius-Mu
 HOL-Number-Theory.Number-Theory
 begin

lemma *fds-totient-times-zeta*:
 $f\text{ds}\ (\lambda n. \text{of-nat}\ (\text{totient}\ n) :: 'a :: \text{comm-semiring-1}) * f\text{ds-zeta} = f\text{ds}\ \text{of-nat}$
 <proof>

lemma *fds-totient-times-zeta'*: $f\text{ds}\ \text{totient} * f\text{ds-zeta} = f\text{ds}\ \text{id}$

```

    <proof>

lemma fds-totient:  $\text{fds } (\lambda n. \text{of-nat } (\text{totient } n)) = \text{fds of-nat} * \text{fds moebius-mu}$ 
    <proof>

lemma totient-conv-moebius-mu:
     $\text{int } (\text{totient } n) = \text{dirichlet-prod moebius-mu int } n$ 
    <proof>

interpretation totient: multiplicative-function totient
    <proof>

lemma even-prime-nat:  $\text{prime } p \implies \text{even } p \implies p = (2::\text{nat})$ 
    <proof>

lemma twopow-dvd-totient:
    fixes  $n :: \text{nat}$ 
    assumes  $n > 0$ 
    defines  $k \equiv \text{card } \{p \in \text{prime-factors } n. \text{odd } p\}$ 
    shows  $2^k \text{ dvd totient } n$ 
    <proof>

lemma totient-conv-moebius-mu':
    assumes  $n > (0::\text{nat})$ 
    shows  $\text{real } (\text{totient } n) = \text{real } n * (\sum d \mid d \text{ dvd } n. \text{moebius-mu } d / \text{real } d)$ 
    <proof>

lemma totient-prime-power-Suc:
    assumes prime  $p$ 
    shows  $\text{totient } (p^{\text{Suc } n}) = p^{\text{Suc } n} - p^n$ 
    <proof>

interpretation totient: multiplicative-function'  $\text{totient } \lambda p k. p^k - p^{(k-1)}$ 
     $\lambda p. p - 1$ 
    <proof>

end

```

7 The Liouville λ function

```

theory Liouville-Lambda
  imports
    HOL-Computational-Algebra.Computational-Algebra
    HOL-Number-Theory.Number-Theory
    Dirichlet-Series
    Multiplicative-Function
    Moebius-Mu
begin

```

definition *liouville-lambda* :: nat \Rightarrow 'a :: comm-ring-1 **where**

liouville-lambda n = (if n = 0 then 0 else (-1) \wedge size (prime-factorization n))

interpretation *liouville-lambda*: completely-multiplicative-function' *liouville-lambda*
 λ -. -1

\langle proof \rangle

lemma *liouville-lambda-prime* [simp]: prime p \implies *liouville-lambda* p = -1

\langle proof \rangle

lemma *liouville-lambda-prime-power* [simp]: prime p \implies *liouville-lambda* (p \wedge k)
= (-1) \wedge k

\langle proof \rangle

lemma *liouville-lambda-squarefree*: squarefree n \implies *liouville-lambda* n = moebius-mu n

\langle proof \rangle

lemma *power-neg-one-If*: (-1) \wedge n = (if even n then 1 else -1 :: 'a :: ring-1)

\langle proof \rangle

lemma *liouville-lambda-power-even*:

n > 0 \implies even m \implies *liouville-lambda* (n \wedge m) = 1

\langle proof \rangle

lemma *liouville-lambda-power-odd*:

odd m \implies *liouville-lambda* (n \wedge m) = *liouville-lambda* n

\langle proof \rangle

lemma *liouville-lambda-power*:

liouville-lambda (n \wedge m) =

(if n = 0 \wedge m > 0 then 0 else if even m then 1 else *liouville-lambda* n)

\langle proof \rangle

interpretation *squarefree*: multiplicative-function'

ind squarefree λ p k. if k > 1 then 0 else 1 λ -. 1

\langle proof \rangle

interpretation *is-nth-power*: multiplicative-function' ind (is-nth-power n)

\langle proof \rangle

interpretation *is-nth-power*: multiplicative-function'

ind (is-nth-power n) λ p k. if n dvd k then 1 else 0 λ -. if n = 1 then 1 else 0

\langle proof \rangle

interpretation *is-square*: multiplicative-function' ind is-square

\langle proof \rangle

interpretation *is-square: multiplicative-function'*
ind is-square $\lambda p k$. if even k then 1 else 0 λ . 0
 ⟨proof⟩

lemma *liouville-lambda-divisors-sum:*
 $(\sum d \mid d \text{ dvd } n. \text{liouville-lambda } d) = \text{ind is-square } n$
 ⟨proof⟩

lemma *fds-liouville-lambda-times-zeta:* $\text{fds liouville-lambda} * \text{fds-zeta} = \text{fds-ind is-square}$
 ⟨proof⟩

lemma *fds-liouville-lambda:* $\text{fds liouville-lambda} = \text{fds-ind is-square} * \text{fds moebius-mu}$
 ⟨proof⟩

lemma *liouville-lambda-altdef:*
 $\text{liouville-lambda } n = (\sum d \mid d^2 \text{ dvd } n. \text{moebius-mu } (n \text{ div } d^2))$
 ⟨proof⟩

lemma *abs-moebius-mu:* $\text{abs } (\text{moebius-mu } n :: 'a :: \text{linordered-idom}) = \text{ind square-free } n$
 ⟨proof⟩

end

8 The divisor functions

theory *Divisor-Count*

imports

Complex-Main

HOL-Number-Theory.Number-Theory

Dirichlet-Series

More-Totient

Moebius-Mu

begin

8.1 The general divisor function

definition *divisor-sigma* $:: 'a :: \text{nat-power} \Rightarrow \text{nat} \Rightarrow 'a$ **where**
 $\text{divisor-sigma } x n = (\sum d \mid d \text{ dvd } n. \text{nat-power } d x)$

lemma *divisor-sigma-0 [simp]:* $\text{divisor-sigma } x 0 = 0$
 ⟨proof⟩

lemma *divisor-sigma-Suc-0 [simp]:* $\text{divisor-sigma } x (\text{Suc } 0) = 1$
 ⟨proof⟩

lemma *divisor-sigma-1* [simp]: $\text{divisor-sigma } x \ 1 = 1$
 ⟨proof⟩

lemma *fds-divisor-sigma*: $\text{fds } (\text{divisor-sigma } x) = \text{fds-zeta} * \text{fds-shift } x \ \text{fds-zeta}$
 ⟨proof⟩

interpretation *divisor-sigma*: *multiplicative-function divisor-sigma* x
 ⟨proof⟩

lemma *divisor-sigma-naive* [code]:
 $\text{divisor-sigma } x \ n = (\text{if } n = 0 \text{ then } 0 \text{ else fold-atLeastAtMost-nat } (\lambda d \text{ acc. if } d \ \text{dvd } n \text{ then nat-power } d \ x + \text{acc else acc}) \ 1 \ n \ 0)$
 ⟨proof⟩

lemma *divisor-sigma-of-nat*: $\text{divisor-sigma } (\text{of-nat } x) \ n = \text{of-nat } (\text{divisor-sigma } x \ n)$
 ⟨proof⟩

lemma *divisor-sigma-prime-power-field*:
fixes $x :: 'a :: \{\text{field}, \text{nat-power}\}$
assumes *prime* p
shows $\text{divisor-sigma } x \ (p \ ^k) =$
 $(\text{if nat-power } p \ x = 1 \text{ then of-nat } (k + 1) \text{ else } (\text{nat-power } p \ x \ ^{\text{Suc } k - 1}) / (\text{nat-power } p \ x - 1))$
 ⟨proof⟩

lemma *divisor-sigma-prime-power-nat*:
assumes *prime* p
shows $\text{divisor-sigma } x \ (p \ ^k) = (\text{if } x = 0 \text{ then Suc } k \text{ else } (p \ ^{(x * \text{Suc } k) - 1}) \ \text{div } (p \ ^x - 1))$
 ⟨proof⟩

interpretation *divisor-sigma-field*:
multiplicative-function' *divisor-sigma* $(x :: 'a :: \{\text{field}, \text{nat-power}\})$
 $\lambda p \ k. \text{if nat-power } p \ x = 1 \text{ then of-nat } (\text{Suc } k) \text{ else } (\text{nat-power } p \ x \ ^{\text{Suc } k - 1}) / (\text{nat-power } p \ x - 1)$
 $\lambda p. \text{nat-power } p \ x + 1$
 ⟨proof⟩

interpretation *divisor-sigma-real*:
multiplicative-function' *divisor-sigma* $(x :: \text{real})$
 $\lambda p \ k. \text{if } x = 0 \text{ then of-nat } (\text{Suc } k) \text{ else } ((\text{real } p \ \text{powr } x) \ ^{\text{Suc } k - 1}) / (\text{real } p \ \text{powr } x - 1)$
 $\lambda p. \text{real } p \ \text{powr } x + 1$
 ⟨proof⟩

interpretation *divisor-sigma-nat*:
multiplicative-function' *divisor-sigma* $(x :: \text{nat})$
 $\lambda p \ k. \text{if } x = 0 \text{ then Suc } k \text{ else } (p \ ^{(\text{Suc } k * x) - 1}) \ \text{div } (p \ ^x - 1)$

$\lambda p. p \wedge x + 1$
<proof>

lemma *divisor-sigma-prime*:
 assumes *prime p*
 shows *divisor-sigma x p = nat-power p x + 1*
<proof>

8.2 The divisor-counting function

definition *divisor-count* :: *nat* \Rightarrow *nat* **where**
 divisor-count n = card {d. d dvd n}

lemma *divisor-count-0* [*simp*]: *divisor-count 0 = 0*
<proof>

lemma *divisor-count-Suc-0* [*simp*]: *divisor-count (Suc 0) = 1*
<proof>

lemma *divisor-sigma-0-left-nat*: *divisor-sigma 0 n = divisor-count n*
<proof>

lemma *divisor-sigma-0-left*: *divisor-sigma 0 n = of-nat (divisor-count n)*
<proof>

lemma *divisor-count-altdef*: *divisor-count n = divisor-sigma 0 n*
<proof>

lemma *divisor-count-naive* [*code*]:
 divisor-count n = (if n = 0 then 0 else
 fold-atLeastAtMost-nat (λd acc. if d dvd n then Suc acc else acc) 1 n 0)
<proof>

interpretation *divisor-count*: *multiplicative-function'* *divisor-count* $\lambda p k. Suc k$
 $\lambda-. 2$
<proof>

lemma *divisor-count-dvd-mono*:
 assumes *a dvd b b \neq 0*
 shows *divisor-count a \leq divisor-count b*
<proof>

8.3 The divisor sum function

definition *divisor-sum* :: *nat* \Rightarrow *nat* **where**
 divisor-sum n = $\sum \{d. d dvd n\}$

lemma *divisor-sum-0* [*simp*]: *divisor-sum 0 = 0*
<proof>

lemma *divisor-sum-Suc-0* [simp]: $\text{divisor-sum } (\text{Suc } 0) = \text{Suc } 0$
(proof)

lemma *divisor-sigma-1-left-nat*: $\text{divisor-sigma } (\text{Suc } 0) n = \text{divisor-sum } n$
(proof)

lemma *divisor-sigma-1-left*: $\text{divisor-sigma } 1 n = \text{of-nat } (\text{divisor-sum } n)$
(proof)

lemma *divisor-sum-altdef*: $\text{divisor-sum } n = \text{divisor-sigma } 1 n$
(proof)

interpretation *divisor-sum*:

multiplicative-function' *divisor-sum* $\lambda p k. (p \wedge \text{Suc } k - 1) \text{div } (p - 1) \lambda p. \text{Suc } p$
(proof)

lemma *divisor-sum-dvd-mono*:

assumes $a \text{ dvd } b \ b \neq 0$

shows $\text{divisor-sum } a \leq \text{divisor-sum } b$

(proof)

lemma *divisor-sum-naive* [code]:

$\text{divisor-sum } n = (\text{if } n = 0 \text{ then } 0 \text{ else}$

$\text{fold-atLeastAtMost-nat } (\lambda d \text{ acc. if } d \text{ dvd } n \text{ then } d + \text{acc} \text{ else } \text{acc}) 1 n 0)$

(proof)

lemma *fds-divisor-count*: $\text{fds divisor-count} = \text{fds-zeta } ^\wedge 2$
(proof)

lemma *fds-shift-zeta-1*: $\text{fds-shift } 1 \text{fds-zeta} = \text{fds of-nat}$
(proof)

lemma *fds-shift-zeta-Suc-0*: $\text{fds-shift } (\text{Suc } 0) \text{fds-zeta} = \text{fds id}$
(proof)

lemma *fds-divisor-sum*: $\text{fds divisor-sum} = \text{fds-zeta} * \text{fds id}$
(proof)

lemma *fds-divisor-sum-eq-totient-times-d*: $\text{fds divisor-sum} = \text{fds totient} * \text{fds divisor-count}$
(proof)

lemma *fds-divisor-sum-times-moebius-mu*:

$\text{fds } (\text{divisor-sigma } (1 :: 'a :: \{\text{nat-power, comm-ring-1}\})) * \text{fds moebius-mu} = \text{fds of-nat}$

(proof)

lemma *inverse-divisor-sigma*:
fixes $a :: 'a :: \{\text{field}, \text{nat-power}\}$
shows $\text{inverse} (\text{fds} (\text{divisor-sigma } a)) = \text{fds-shift } a (\text{fds } \text{moebius-mu}) * \text{fds } \text{moebius-mu}$
 $\langle \text{proof} \rangle$

end

9 Summatory arithmetic functions

theory *Arithmetic-Summatory*

imports

More-Totient
Moebius-Mu
Liouville-Lambda
Divisor-Count
Dirichlet-Series

begin

9.1 Definition

definition *sum-upto* :: $(\text{nat} \Rightarrow 'a :: \text{comm-monoid-add}) \Rightarrow \text{real} \Rightarrow 'a$ **where**
 $\text{sum-upto } f \ x = (\sum i \mid 0 < i \wedge \text{real } i \leq x. f \ i)$

lemma *sum-upto-altdef*: $\text{sum-upto } f \ x = (\sum i \in \{0 <.. \text{nat } \lfloor x \rfloor\}. f \ i)$
 $\langle \text{proof} \rangle$

lemma *sum-upto-0* [*simp*]: $\text{sum-upto } f \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma *sum-upto-cong* [*cong*]:
 $(\bigwedge n. n > 0 \implies f \ n = f' \ n) \implies n = n' \implies \text{sum-upto } f \ n = \text{sum-upto } f' \ n'$
 $\langle \text{proof} \rangle$

lemma *finite-Nats-le-real* [*simp,intro*]: $\text{finite } \{n. 0 < n \wedge \text{real } n \leq x\}$
 $\langle \text{proof} \rangle$

lemma *sum-upto-ind*: $\text{sum-upto} (\text{ind } P) \ x = \text{of-nat} (\text{card } \{n. n > 0 \wedge \text{real } n \leq x \wedge P \ n\})$
 $\langle \text{proof} \rangle$

lemma *sum-upto-sum-divisors*:
 $\text{sum-upto} (\lambda n. \sum d \mid d \text{ dvd } n. f \ n \ d) \ x = \text{sum-upto} (\lambda k. \text{sum-upto} (\lambda d. f \ (d * k) \ k) \ (x / k)) \ x$
 $\langle \text{proof} \rangle$

lemma *sum-upto-dirichlet-prod*:
 $\text{sum-upto} (\text{dirichlet-prod } f \ g) \ x = \text{sum-upto} (\lambda d. f \ d * \text{sum-upto } g \ (x / \text{real } d)) \ x$

<proof>

lemma *sum-upto-real*:

assumes $x \geq 0$

shows $\text{sum-upto real } x = \text{of-int (floor } x) * (\text{of-int (floor } x) + 1) / 2$

<proof>

lemma *summable-imp-convergent-sum-upto*:

assumes *summable* ($f :: \text{nat} \Rightarrow 'a :: \text{real-normed-vector}$)

obtains c **where** ($\text{sum-upto } f \longrightarrow c$) *at-top*

<proof>

9.2 The Hyperbola method

lemma *hyperbola-method-semiring*:

fixes $f g :: \text{nat} \Rightarrow 'a :: \text{comm-semiring-0}$

assumes $A \geq 0$ **and** $B \geq 0$ **and** $A * B = x$

shows $\text{sum-upto (dirichlet-prod } f g) x + \text{sum-upto } f A * \text{sum-upto } g B =$
 $\text{sum-upto } (\lambda n. f n * \text{sum-upto } g (x / \text{real } n)) A +$
 $\text{sum-upto } (\lambda n. \text{sum-upto } f (x / \text{real } n) * g n) B$

<proof>

lemma *hyperbola-method-semiring-sqrt*:

fixes $f g :: \text{nat} \Rightarrow 'a :: \text{comm-semiring-0}$

assumes $x \geq 0$

shows $\text{sum-upto (dirichlet-prod } f g) x + \text{sum-upto } f (\text{sqrt } x) * \text{sum-upto } g (\text{sqrt } x) =$

$\text{sum-upto } (\lambda n. f n * \text{sum-upto } g (x / \text{real } n)) (\text{sqrt } x) +$
 $\text{sum-upto } (\lambda n. \text{sum-upto } f (x / \text{real } n) * g n) (\text{sqrt } x)$

<proof>

lemma *hyperbola-method*:

fixes $f g :: \text{nat} \Rightarrow 'a :: \text{comm-ring}$

assumes $A \geq 0$ $B \geq 0$ $A * B = x$

shows $\text{sum-upto (dirichlet-prod } f g) x =$
 $\text{sum-upto } (\lambda n. f n * \text{sum-upto } g (x / \text{real } n)) A +$
 $\text{sum-upto } (\lambda n. \text{sum-upto } f (x / \text{real } n) * g n) B -$
 $\text{sum-upto } f A * \text{sum-upto } g B$

<proof>

lemma *hyperbola-method-sqrt*:

fixes $f g :: \text{nat} \Rightarrow 'a :: \text{comm-ring}$

assumes $x \geq 0$

shows $\text{sum-upto (dirichlet-prod } f g) x =$
 $\text{sum-upto } (\lambda n. f n * \text{sum-upto } g (x / \text{real } n)) (\text{sqrt } x) +$
 $\text{sum-upto } (\lambda n. \text{sum-upto } f (x / \text{real } n) * g n) (\text{sqrt } x) -$
 $\text{sum-upto } f (\text{sqrt } x) * \text{sum-upto } g (\text{sqrt } x)$

<proof>

end

10 Partial summation

theory *Partial-Summation*

imports

HOL-Analysis.Analysis

Arithmetic-Summatory

begin

lemma *finite-vimage-real-of-nat-greaterThanAtMost*: *finite* (real - ' {y<..x})
(*proof*)

context

fixes *a* :: nat \Rightarrow 'a :: {banach, real-normed-algebra}

fixes *f f'* :: real \Rightarrow 'a

fixes *A*

fixes *X* :: real set

fixes *x y* :: real

defines *A* \equiv *sum-upto a*

assumes *fin*: *finite X*

assumes *xy*: $0 \leq y < x$

assumes *deriv*: $\bigwedge z. z \in \{y..x\} - X \implies (f \text{ has-vector-derivative } f' z) (at z)$

assumes *cont-f*: *continuous-on* {y..x} *f*

begin

lemma *partial-summation-strong*:

$((\lambda t. A t * f' t) \text{ has-integral}$

$(A x * f x - A y * f y - (\sum n \in \text{real} - ' \{y<..x\}. a n * f n))) \{y..x\}$

(*proof*)

lemma *partial-summation-integrable-strong*:

$(\lambda t. A t * f' t) \text{ integrable-on } \{y..x\}$

and *partial-summation-strong'*:

$(\sum n \in \text{real} - ' \{y<..x\}. a n * f n) =$

$A x * f x - A y * f y - \text{integral } \{y..x\} (\lambda t. A t * f' t)$

(*proof*)

end

context

fixes *a* :: nat \Rightarrow 'a :: {banach, real-normed-algebra}

fixes *f f'* :: real \Rightarrow 'a

fixes *A*

fixes *X* :: real set

fixes *x* :: real

defines *A* \equiv *sum-upto a*

assumes *fin*: *finite X*

assumes $x: x > 0$
assumes $deriv: \bigwedge z. z \in \{0..x\} - X \implies (f \text{ has-vector-derivative } f' z) \text{ (at } z)$
assumes $cont-f: \text{continuous-on } \{0..x\} f$
begin

lemma *partial-summation-sum-upto-strong*:
 $((\lambda t. A t * f' t) \text{ has-integral } (A x * f x - \text{sum-upto } (\lambda n. a n * f n) x)) \{0..x\}$
 $\langle \text{proof} \rangle$

lemma *partial-summation-integrable-sum-upto-strong*:
 $(\lambda t. A t * f' t) \text{ integrable-on } \{0..x\}$
and *partial-summation-sum-upto-strong'*:
 $\text{sum-upto } (\lambda n. a n * f n) x =$
 $A x * f x - \text{integral } \{0..x\} (\lambda t. A t * f' t)$
 $\langle \text{proof} \rangle$

end

end

11 Euler product expansions

theory *Euler-Products*

imports

HOL-Analysis.Analysis

Multiplicative-Function

begin

Conflicting notation from *HOL-Analysis.Infinite-Sum*

no-notation *Infinite-Sum.abs-summable-on* (**infixr** $\langle \text{abs}'\text{-summable}'\text{-on} \rangle$ 46)

lemma *prime-factors-power-subset*:
 $\text{prime-factors } (x \wedge n) \subseteq \text{prime-factors } x$
 $\langle \text{proof} \rangle$

lemma *prime-power-product-in-Pi*:
 $(\lambda g. \prod p \in \{p. p \leq (n::\text{nat}) \wedge \text{prime } p\}. p \wedge g p)$
 $\in (\{p. p \leq n \wedge \text{prime } p\} \rightarrow_E \text{UNIV}) \rightarrow$
 $\{m. 0 < m \wedge \text{prime-factors } m \subseteq \{..n\}\}$
 $\langle \text{proof} \rangle$

lemma *inj-prime-power*: $\text{inj-on } (\lambda x. \text{fst } x \wedge \text{snd } x :: \text{nat}) (\{a. \text{prime } a\} \times \{0<..\})$
 $\langle \text{proof} \rangle$

lemma *bij-betw-prime-powers*:
 $\text{bij-betw } (\lambda g. \prod p \in \{p. p \leq n \wedge \text{prime } p\}. p \wedge g p) (\{p. p \leq n \wedge \text{prime } p\} \rightarrow_E \text{UNIV})$
 $\{m. 0 < m \wedge \text{prime-factors } m \subseteq \{..(n::\text{nat})\}\}$
 $\langle \text{proof} \rangle$

```

lemma
  fixes  $f :: nat \Rightarrow 'a :: \{real-normed-field, banach, second-countable-topology\}$ 
  assumes summable: summable ( $\lambda n. norm (f n)$ )
  assumes multiplicative-function  $f$ 
  shows abs-convergent-euler-product:
    abs-convergent-prod ( $\lambda p. \text{if prime } p \text{ then } \sum n. f (p \wedge n) \text{ else } 1$ )
    and euler-product-LIMSEQ:
      ( $\lambda n. (\prod p \leq n. \text{if prime } p \text{ then } \sum n. f (p \wedge n) \text{ else } 1)$ )  $\longrightarrow$  ( $\sum n. f n$ )
  <proof>

```

```

lemma
  fixes  $f :: nat \Rightarrow 'a :: \{real-normed-field, banach, second-countable-topology\}$ 
  assumes summable: summable ( $\lambda n. norm (f n)$ )
  assumes completely-multiplicative-function  $f$ 
  shows abs-convergent-euler-product':
    abs-convergent-prod ( $\lambda p. \text{if prime } p \text{ then inverse } (1 - f p) \text{ else } 1$ )
    and completely-multiplicative-summable-norm:
       $\bigwedge p. \text{prime } p \implies norm (f p) < 1$ 
    and euler-product-LIMSEQ':
      ( $\lambda n. (\prod p \leq n. \text{if prime } p \text{ then inverse } (1 - f p) \text{ else } 1)$ )  $\longrightarrow$  ( $\sum n. f$ 
   $n$ )
  <proof>

```

end

12 Analytic properties of Dirichlet series

theory *Dirichlet-Series-Analysis*

imports

HOL-Complex-Analysis.Complex-Analysis
HOL-Library.Going-To-Filter
HOL-Real-Asymp.Real-Asymp
Dirichlet-Series
Moebius-Mu
Partial-Summation
Euler-Products

begin

Conflicting notation from *HOL-Analysis.Infinite-Sum*

no-notation *Infinite-Sum.abs-summable-on* (**infixr** $\langle \text{abs}'\text{-summable}'\text{-on} \rangle$ 46)

The following illustrates a concept we will need later on: A property holds for f going to F if we can find e. g. a sequence that tends to F and whose elements eventually satisfy P .

```

lemma frequently-going-toI:
  assumes filterlim ( $\lambda n. f (g n)$ )  $F G$ 
  assumes eventually ( $\lambda n. P (g n)$ )  $G$ 

```

assumes *eventually* $(\lambda n. g\ n \in A)$ G
assumes $G \neq \text{bot}$
shows *frequently* P (*f going-to F within A*)
 $\langle \text{proof} \rangle$

lemma *frequently-filtercomapI*:
assumes *filterlim* $(\lambda n. f\ (g\ n))$ $F\ G$
assumes *eventually* $(\lambda n. P\ (g\ n))$ G
assumes $G \neq \text{bot}$
shows *frequently* P (*filtercomap f F*)
 $\langle \text{proof} \rangle$

lemma *frequently-going-to-at-topE*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes *frequently* P (*f going-to at-top*)
obtains g **where** $\bigwedge n. P\ (g\ n)$ **and** *filterlim* $(\lambda n. f\ (g\ n))$ *at-top sequentially*
 $\langle \text{proof} \rangle$

Apostol often uses statements like ‘ $P(s_k)$ for all k in an infinite sequence s_k such that $\mathfrak{R}(s_k) \rightarrow \infty$ as $k \rightarrow \infty$ ’.

Instead, we write *frequently P (Re going-to at-top)*. This lemma shows that our statement is equivalent to his.

lemma *frequently-going-to-at-top-iff*:
frequently P (f going-to (at-top :: real filter)) \longleftrightarrow
 $(\exists g. \forall n. P\ (g\ n) \wedge \text{filterlim}\ (\lambda n. f\ (g\ n))\ \text{at-top sequentially})$
 $\langle \text{proof} \rangle$

lemma *surj-bullet-1*: *surj* $(\lambda s :: 'a :: \{\text{real-normed-algebra-1}, \text{real-inner}\}. s \cdot 1)$
 $\langle \text{proof} \rangle$

lemma *bullet-1-going-to-at-top-neq-bot [simp]*:
 $((\lambda s :: 'a :: \{\text{real-normed-algebra-1}, \text{real-inner}\}. s \cdot 1)\ \text{going-to at-top}) \neq \text{bot}$
 $\langle \text{proof} \rangle$

lemma *fds-abs-converges-altdef*:
fds-abs-converges f s $\longleftrightarrow (\lambda n. \text{fds-nth}\ f\ n / \text{nat-power}\ n\ s)\ \text{abs-summable-on}\ \{1..\}$
 $\langle \text{proof} \rangle$

lemma *fds-abs-converges-altdef'*:
fds-abs-converges f s $\longleftrightarrow (\lambda n. \text{fds-nth}\ f\ n / \text{nat-power}\ n\ s)\ \text{abs-summable-on}\ \text{UNIV}$
 $\langle \text{proof} \rangle$

lemma *eval-fds-altdef*:
assumes *fds-abs-converges f s*
shows $\text{eval-fds}\ f\ s = (\sum_a n. \text{fds-nth}\ f\ n / \text{nat-power}\ n\ s)$

<proof>

lemma *multiplicative-function-divide-nat-power:*

fixes $f :: \text{nat} \Rightarrow 'a :: \{\text{nat-power}, \text{field}\}$

assumes *multiplicative-function* f

shows *multiplicative-function* $(\lambda n. f\ n / \text{nat-power}\ n\ s)$

<proof>

lemma *completely-multiplicative-function-divide-nat-power:*

fixes $f :: \text{nat} \Rightarrow 'a :: \{\text{nat-power}, \text{field}\}$

assumes *completely-multiplicative-function* f

shows *completely-multiplicative-function* $(\lambda n. f\ n / \text{nat-power}\ n\ s)$

<proof>

12.1 Convergence and absolute convergence

class *nat-power-normed-field* = *nat-power-field* + *real-normed-field* + *real-inner* + *real-algebra-1* +

fixes *real-power* :: $\text{real} \Rightarrow 'a \Rightarrow 'a$

assumes *real-power-nat-power*: $n > 0 \implies \text{real-power}\ (\text{real}\ n)\ c = \text{nat-power}\ n\ c$

assumes *real-power-1-right-aux*: $d > 0 \implies \text{real-power}\ d\ 1 = d *_{\mathbb{R}} 1$

assumes *real-power-add*: $d > 0 \implies \text{real-power}\ d\ (a + b) = \text{real-power}\ d\ a * \text{real-power}\ d\ b$

assumes *real-power-nonzero* [*simp*]: $d > 0 \implies \text{real-power}\ d\ a \neq 0$

assumes *norm-real-power*: $x > 0 \implies \text{norm}\ (\text{real-power}\ x\ c) = x\ \text{powr}\ (c \cdot 1)$

assumes *nat-power-of-real-aux*: $\text{nat-power}\ n\ (x *_{\mathbb{R}} 1) = ((\text{real}\ n\ \text{powr}\ x) *_{\mathbb{R}} 1)$

assumes *has-field-derivative-nat-power-aux*:

$\bigwedge x::'a. n > 0 \implies \text{LIM}\ y\ \text{inf-class.inf}$

$(\text{Inf}\ (\text{principal}\ ' \{S. \text{open}\ } S \wedge x \in S))\ (\text{principal}\ (UNIV - \{x\}))$

$(\text{nat-power}\ n\ y - \text{nat-power}\ n\ x - \ln\ (\text{real}\ n) *_{\mathbb{R}} \text{nat-power}\ n\ x * (y - x)) /_{\mathbb{R}}$

$\text{norm}\ (y - x) := \text{Inf}\ (\text{principal}\ ' \{S. \text{open}\ } S \wedge 0 \in S)$

assumes *has-vector-derivative-real-power-aux*:

$x > 0 \implies \text{filterlim}\ (\lambda y. (\text{real-power}\ y\ c - \text{real-power}\ x\ c - \text{real-power}\ x\ (c - 1)) -$

$(y - x) *_{\mathbb{R}} (c * \text{real-power}\ x\ (c - 1))) /_{\mathbb{R}}$

$\text{norm}\ (y - x)\ (\text{INF}\ S \in \{S. \text{open}\ } S \wedge 0 \in S). \text{principal}\ S)\ (\text{at}\ x)$

assumes *norm-nat-power*: $n > 0 \implies \text{norm}\ (\text{nat-power}\ n\ y) = \text{real}\ n\ \text{powr}\ (y \cdot 1)$

begin

lemma *real-power-diff*: $d > 0 \implies \text{real-power}\ d\ (a - b) = \text{real-power}\ d\ a / \text{real-power}\ d\ b$

<proof>

end

lemma *real-power-1-right* [*simp*]: $d > 0 \implies \text{real-power}\ d\ 1 = \text{of-real}\ d$

<proof>

lemma *has-vector-derivative-real-power* [*derivative-intros*]:
 $x > 0 \implies ((\lambda y. \text{real-power } y \ c) \text{ has-vector-derivative } c * \text{real-power } x \ (c - 1))$
(at x within A)
⟨*proof*⟩

lemma *has-field-derivative-nat-power* [*derivative-intros*]:
 $n > 0 \implies ((\lambda y. \text{nat-power } n \ y) \text{ has-field-derivative } \ln \ (\text{real } n) *_{\mathbb{R}} \text{nat-power } n \ x)$
(at $(x :: 'a :: \text{nat-power-normed-field})$ within A)
⟨*proof*⟩

lemma *continuous-on-real-power* [*continuous-intros*]:
 $A \subseteq \{0 < ..\} \implies \text{continuous-on } A \ (\lambda x. \text{real-power } x \ s)$
⟨*proof*⟩

instantiation *real* :: *nat-power-normed-field*
begin

definition *real-power-real* :: *real* \Rightarrow *real* \Rightarrow *real* **where**
[*simp*]: *real-power-real* = (*powr*)

instance ⟨*proof*⟩

end

instantiation *complex* :: *nat-power-normed-field*
begin

definition *nat-power-complex* :: *nat* \Rightarrow *complex* \Rightarrow *complex* **where**
[*simp*]: *nat-power-complex* $n \ z = \text{of-nat } n \ \text{powr } z$

definition *real-power-complex* :: *real* \Rightarrow *complex* \Rightarrow *complex* **where**
[*simp*]: *real-power-complex* = $(\lambda x \ y. \text{of-real } x \ \text{powr } y)$

instance ⟨*proof*⟩

end

lemma *nat-power-of-real* [*simp*]:
 $\text{nat-power } n \ (\text{of-real } x :: 'a :: \text{nat-power-normed-field}) = \text{of-real } (\text{real } n \ \text{powr } x)$
⟨*proof*⟩

lemma *fds-abs-converges-of-real* [*simp*]:
fds-abs-converges (*fds-of-real* f)
 $(\text{of-real } s :: 'a :: \{\text{nat-power-normed-field}, \text{banach}\}) \iff \text{fds-abs-converges } f \ s$
⟨*proof*⟩

lemma *eval-fds-of-real* [*simp*]:

assumes $fds\text{-converges } f s$
shows $eval\text{-fds } (fds\text{-of-real } f) (of\text{-real } s :: 'a :: \{nat\text{-power-normed-field}, banach\})$
 $=$
 $of\text{-real } (eval\text{-fds } f s)$
 $\langle proof \rangle$

lemma $fds\text{-abs-summable-zeta-iff [simp]}$:
fixes $s :: 'a :: \{banach, nat\text{-power-normed-field}\}$
shows $fds\text{-abs-converges } fds\text{-zeta } s \longleftrightarrow s \cdot 1 > (1 :: real)$
 $\langle proof \rangle$

lemma $fds\text{-abs-summable-zeta}$:
 $(s :: 'a :: \{banach, nat\text{-power-normed-field}\}) \cdot 1 > 1 \implies fds\text{-abs-converges } fds\text{-zeta } s$
 $\langle proof \rangle$

lemma $fds\text{-abs-converges-moebius-mu}$:
fixes $s :: 'a :: \{banach, nat\text{-power-normed-field}\}$
assumes $s \cdot 1 > 1$
shows $fds\text{-abs-converges } (fds\text{ moebius-mu}) s$
 $\langle proof \rangle$

definition $conv\text{-abscissa}$
 $:: 'a :: \{nat\text{-power}, banach, real\text{-normed-field}, real\text{-inner}\} fds \Rightarrow ereal \textbf{ where}$
 $conv\text{-abscissa } f = (INF s \in \{s. fds\text{-converges } f s\}. ereal (s \cdot 1))$

definition $abs\text{-conv-abscissa}$
 $:: 'a :: \{nat\text{-power}, banach, real\text{-normed-field}, real\text{-inner}\} fds \Rightarrow ereal \textbf{ where}$
 $abs\text{-conv-abscissa } f = (INF s \in \{s. fds\text{-abs-converges } f s\}. ereal (s \cdot 1))$

lemma $conv\text{-abscissa-mono}$:
assumes $\bigwedge s. fds\text{-converges } g s \implies fds\text{-converges } f s$
shows $conv\text{-abscissa } f \leq conv\text{-abscissa } g$
 $\langle proof \rangle$

lemma $abs\text{-conv-abscissa-mono}$:
assumes $\bigwedge s. fds\text{-abs-converges } g s \implies fds\text{-abs-converges } f s$
shows $abs\text{-conv-abscissa } f \leq abs\text{-conv-abscissa } g$
 $\langle proof \rangle$

class $dirichlet\text{-series} = euclidean\text{-space} + real\text{-normed-field} + nat\text{-power-normed-field}$
 $+$
assumes $one\text{-in-Basis}: 1 \in Basis$

instance $real :: dirichlet\text{-series} \langle proof \rangle$
instance $complex :: dirichlet\text{-series} \langle proof \rangle$

context

assumes *SORT-CONSTRAINT*('a :: *dirichlet-series*)

begin

lemma *fds-abs-converges-Re-le*:

fixes $f :: 'a \text{ fds}$

assumes $\text{fds-abs-converges } f \ z \ z \cdot 1 \leq z' \cdot 1$

shows $\text{fds-abs-converges } f \ z'$

<proof>

lemma *fds-abs-converges*:

assumes $s \cdot 1 > \text{abs-conv-abscissa } (f :: 'a \text{ fds})$

shows $\text{fds-abs-converges } f \ s$

<proof>

lemma *fds-abs-diverges*:

assumes $s \cdot 1 < \text{abs-conv-abscissa } (f :: 'a \text{ fds})$

shows $\neg \text{fds-abs-converges } f \ s$

<proof>

lemma *uniformly-Cauchy-eval-fds-aux*:

fixes $s0 :: 'a :: \text{dirichlet-series}$

assumes $\text{bounded: } B \text{seq } (\lambda n. \sum_{k \leq n}. \text{fds-nth } f \ k / \text{nat-power } k \ s0)$

assumes $B: \text{compact } B \ \bigwedge z. z \in B \implies z \cdot 1 > s0 \cdot 1$

shows $\text{uniformly-Cauchy-on } B \ (\lambda N \ z. \sum_{n \leq N}. \text{fds-nth } f \ n / \text{nat-power } n \ z)$

<proof>

lemma *uniformly-convergent-eval-fds-aux*:

assumes $B \text{seq } (\lambda n. \sum_{k \leq n}. \text{fds-nth } f \ k / \text{nat-power } k \ (s0 :: 'a))$

assumes $B: \text{compact } B \ \bigwedge z. z \in B \implies z \cdot 1 > s0 \cdot 1$

shows $\text{uniformly-convergent-on } B \ (\lambda N \ z. \sum_{n \leq N}. \text{fds-nth } f \ n / \text{nat-power } n \ z)$

<proof>

lemma *uniformly-convergent-eval-fds-aux'*:

assumes $\text{conv: } \text{fds-converges } f \ (s0 :: 'a)$

assumes $B: \text{compact } B \ \bigwedge z. z \in B \implies z \cdot 1 > s0 \cdot 1$

shows $\text{uniformly-convergent-on } B \ (\lambda N \ z. \sum_{n \leq N}. \text{fds-nth } f \ n / \text{nat-power } n \ z)$

<proof>

lemma *bounded-partial-sums-imp-fps-converges*:

fixes $s0 :: 'a :: \text{dirichlet-series}$

assumes $B \text{seq } (\lambda n. \sum_{k \leq n}. \text{fds-nth } f \ k / \text{nat-power } k \ s0)$ **and** $s \cdot 1 > s0 \cdot 1$

shows $\text{fds-converges } f \ s$

<proof>

theorem *fds-converges-Re-le*:

assumes $\text{fds-converges } f \ (s0 :: 'a) \ s \cdot 1 > s0 \cdot 1$

shows $\text{fds-converges } f \ s$
 $\langle \text{proof} \rangle$

lemma fds-converges :
assumes $s \cdot 1 > \text{conv-abscissa } (f :: 'a \ \text{fds})$
shows $\text{fds-converges } f \ s$
 $\langle \text{proof} \rangle$

lemma fds-diverges :
assumes $s \cdot 1 < \text{conv-abscissa } (f :: 'a \ \text{fds})$
shows $\neg \text{fds-converges } f \ s$
 $\langle \text{proof} \rangle$

theorem $\text{fds-converges-imp-abs-converges}$:
assumes $\text{fds-converges } (f :: 'a \ \text{fds}) \ s \ s' \cdot 1 > s \cdot 1 + 1$
shows $\text{fds-abs-converges } f \ s'$
 $\langle \text{proof} \rangle$

lemma $\text{conv-le-abs-conv-abscissa}$: $\text{conv-abscissa } f \leq \text{abs-conv-abscissa } f$
 $\langle \text{proof} \rangle$

lemma $\text{conv-abscissa-PInf-iff}$: $\text{conv-abscissa } f = \infty \longleftrightarrow (\forall s. \neg \text{fds-converges } f \ s)$
 $\langle \text{proof} \rangle$

lemma $\text{conv-abscissa-PInfI}$ [intro]: $(\bigwedge s. \neg \text{fds-converges } f \ s) \Longrightarrow \text{conv-abscissa } f = \infty$
 $\langle \text{proof} \rangle$

lemma $\text{conv-abscissa-MInf-iff}$: $\text{conv-abscissa } (f :: 'a \ \text{fds}) = -\infty \longleftrightarrow (\forall s. \text{fds-converges } f \ s)$
 $\langle \text{proof} \rangle$

lemma $\text{conv-abscissa-MInfI}$ [intro]: $(\bigwedge s. \text{fds-converges } (f :: 'a \ \text{fds}) \ s) \Longrightarrow \text{conv-abscissa } f = -\infty$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-abscissa-PInf-iff}$: $\text{abs-conv-abscissa } f = \infty \longleftrightarrow (\forall s. \neg \text{fds-abs-converges } f \ s)$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-abscissa-PInfI}$ [intro]: $(\bigwedge s. \neg \text{fds-converges } f \ s) \Longrightarrow \text{abs-conv-abscissa } f = \infty$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-abscissa-MInf-iff}$:
 $\text{abs-conv-abscissa } (f :: 'a \ \text{fds}) = -\infty \longleftrightarrow (\forall s. \text{fds-abs-converges } f \ s)$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-abscissa-MInfI}$ [intro]:

$(\bigwedge s. \text{fds-abs-converges } (f :: 'a \text{ fds}) s) \implies \text{abs-conv-abscissa } f = -\infty$
 ⟨proof⟩

lemma *conv-abscissa-geI*:

assumes $\bigwedge c'. \text{ereal } c' < c \implies \exists s. s \cdot 1 = c' \wedge \neg \text{fds-converges } f s$
shows $\text{conv-abscissa } (f :: 'a \text{ fds}) \geq c$

⟨proof⟩

lemma *conv-abscissa-leI*:

assumes $\bigwedge c'. \text{ereal } c' > c \implies \exists s. s \cdot 1 = c' \wedge \text{fds-converges } f s$
shows $\text{conv-abscissa } (f :: 'a \text{ fds}) \leq c$

⟨proof⟩

lemma *abs-conv-abscissa-geI*:

assumes $\bigwedge c'. \text{ereal } c' < c \implies \exists s. s \cdot 1 = c' \wedge \neg \text{fds-abs-converges } f s$
shows $\text{abs-conv-abscissa } (f :: 'a \text{ fds}) \geq c$

⟨proof⟩

lemma *abs-conv-abscissa-leI*:

assumes $\bigwedge c'. \text{ereal } c' > c \implies \exists s. s \cdot 1 = c' \wedge \text{fds-abs-converges } f s$
shows $\text{abs-conv-abscissa } (f :: 'a \text{ fds}) \leq c$

⟨proof⟩

lemma *conv-abscissa-leI-weak*:

assumes $\bigwedge x. \text{ereal } x > d \implies \text{fds-converges } f \text{ (of-real } x)$
shows $\text{conv-abscissa } (f :: 'a \text{ fds}) \leq d$

⟨proof⟩

lemma *abs-conv-abscissa-leI-weak*:

assumes $\bigwedge x. \text{ereal } x > d \implies \text{fds-abs-converges } f \text{ (of-real } x)$
shows $\text{abs-conv-abscissa } (f :: 'a \text{ fds}) \leq d$

⟨proof⟩

lemma *conv-abscissa-truncate [simp]*:

$\text{conv-abscissa } (\text{fds-truncate } m \text{ (} f :: 'a \text{ fds)}) = -\infty$

⟨proof⟩

lemma *abs-conv-abscissa-truncate [simp]*:

$\text{abs-conv-abscissa } (\text{fds-truncate } m \text{ (} f :: 'a \text{ fds)}) = -\infty$

⟨proof⟩

theorem *abs-conv-le-conv-abscissa-plus-1*: $\text{abs-conv-abscissa } (f :: 'a \text{ fds}) \leq \text{conv-abscissa}$

$f + 1$

⟨proof⟩

lemma *uniformly-convergent-eval-fds*:

assumes $B: \text{compact } B \bigwedge z. z \in B \implies z \cdot 1 > \text{conv-abscissa } (f :: 'a \text{ fds})$

shows *uniformly-convergent-on B* $(\lambda N z. \sum n \leq N. \text{fds-nth } f \ n / \text{nat-power } n \ z)$
 ⟨*proof*⟩

corollary *uniformly-convergent-eval-fds'*:

assumes *B: compact B* $\bigwedge z. z \in B \implies z \cdot 1 > \text{conv-abscissa } (f :: 'a \ \text{fds})$

shows *uniformly-convergent-on B* $(\lambda N z. \sum n < N. \text{fds-nth } f \ n / \text{nat-power } n \ z)$
 ⟨*proof*⟩

12.2 Derivative of a Dirichlet series

lemma *fds-converges-deriv-aux*:

assumes *conv: fds-converges f* $(s0 :: 'a)$ **and** *gt: s · 1 > s0 · 1*

shows *fds-converges (fds-deriv f) s*
 ⟨*proof*⟩

theorem

assumes *s · 1 > conv-abscissa (f :: 'a fds)*

shows *fds-converges-deriv: fds-converges (fds-deriv f) s*

and *has-field-derivative-eval-fds [derivative-intros]:*

(eval-fds f has-field-derivative eval-fds (fds-deriv f) s) (at s within A)

⟨*proof*⟩

lemmas *has-field-derivative-eval-fds' [derivative-intros] =*
DERIV-chain2[OF has-field-derivative-eval-fds]

lemma *continuous-eval-fds [continuous-intros]:*

assumes *s · 1 > conv-abscissa f*

shows *continuous (at s within A) (eval-fds (f :: 'a :: dirichlet-series fds))*
 ⟨*proof*⟩

lemma *continuous-eval-fds' [continuous-intros]:*

fixes *f :: 'a :: dirichlet-series fds*

assumes *continuous (at s within A) g g s · 1 > conv-abscissa f*

shows *continuous (at s within A) ($\lambda x. \text{eval-fds } f \ (g \ x)$)*
 ⟨*proof*⟩

lemma *continuous-on-eval-fds [continuous-intros]:*

fixes *f :: 'a :: dirichlet-series fds*

assumes *A ⊆ {s. s · 1 > conv-abscissa f}*

shows *continuous-on A (eval-fds f)*
 ⟨*proof*⟩

lemma *continuous-on-eval-fds' [continuous-intros]:*

fixes *f :: 'a :: dirichlet-series fds*

assumes *continuous-on A g g ' A ⊆ {s. s · 1 > conv-abscissa f}*

shows *continuous-on A ($\lambda x. \text{eval-fds } f \ (g \ x)$)*
 ⟨*proof*⟩

lemma *conv-abscissa-deriv-le:*

fixes $f :: 'a \text{ fds}$
shows $\text{conv-absclissa } (\text{fds-deriv } f) \leq \text{conv-absclissa } f$
 $\langle \text{proof} \rangle$

lemma *abs-conv-absclissa-integral*:
fixes $f :: 'a \text{ fds}$
shows $\text{abs-conv-absclissa } (\text{fds-integral } a \ f) = \text{abs-conv-absclissa } f$
 $\langle \text{proof} \rangle$

lemma *abs-conv-absclissa-ln*:
 $\text{abs-conv-absclissa } (\text{fds-ln } l \ (f :: 'a :: \text{dirichlet-series } \text{fds})) =$
 $\text{abs-conv-absclissa } (\text{fds-deriv } f / f)$
 $\langle \text{proof} \rangle$

lemma *abs-conv-absclissa-deriv*:
fixes $f :: 'a \text{ fds}$
shows $\text{abs-conv-absclissa } (\text{fds-deriv } f) = \text{abs-conv-absclissa } f$
 $\langle \text{proof} \rangle$

lemma *abs-conv-absclissa-higher-deriv*:
 $\text{abs-conv-absclissa } ((\text{fds-deriv } \hat{\sim} n) \ f) = \text{abs-conv-absclissa } (f :: 'a :: \text{dirichlet-series } \text{fds})$
 $\langle \text{proof} \rangle$

lemma *conv-absclissa-higher-deriv-le*:
 $\text{conv-absclissa } ((\text{fds-deriv } \hat{\sim} n) \ f) \leq \text{conv-absclissa } (f :: 'a :: \text{dirichlet-series } \text{fds})$
 $\langle \text{proof} \rangle$

lemma *abs-conv-absclissa-restrict*:
 $\text{abs-conv-absclissa } (\text{fds-subseries } P \ f) \leq \text{abs-conv-absclissa } f$
 $\langle \text{proof} \rangle$

lemma *eval-fds-deriv*:
fixes $f :: 'a \text{ fds}$
assumes $s \cdot 1 > \text{conv-absclissa } f$
shows $\text{eval-fds } (\text{fds-deriv } f) \ s = \text{deriv } (\text{eval-fds } f) \ s$
 $\langle \text{proof} \rangle$

lemma *eval-fds-higher-deriv*:
assumes $(s :: 'a :: \text{dirichlet-series}) \cdot 1 > \text{conv-absclissa } f$
shows $\text{eval-fds } ((\text{fds-deriv } \hat{\sim} n) \ f) \ s = (\text{deriv } \hat{\sim} n) \ (\text{eval-fds } f) \ s$
 $\langle \text{proof} \rangle$

end

12.3 Multiplication of two series

lemma
fixes $f \ g :: \text{nat} \Rightarrow 'a :: \{\text{banach, real-normed-field, second-countable-topology},$

nat-power}
fixes $s :: 'a$
assumes [*simp*]: $f\ 0 = 0\ g\ 0 = 0$
assumes *summable*: $\text{summable } (\lambda n. \text{norm } (f\ n / \text{nat-power } n\ s))$
 $\text{summable } (\lambda n. \text{norm } (g\ n / \text{nat-power } n\ s))$
shows *summable-dirichlet-prod*: $\text{summable } (\lambda n. \text{norm } (\text{dirichlet-prod } f\ g\ n / \text{nat-power } n\ s))$
and *suminf-dirichlet-prod*:
 $(\sum n. \text{dirichlet-prod } f\ g\ n / \text{nat-power } n\ s) =$
 $(\sum n. f\ n / \text{nat-power } n\ s) * (\sum n. g\ n / \text{nat-power } n\ s)$
<proof>

lemma
fixes $f\ g :: \text{nat} \Rightarrow \text{real}$
fixes $s :: \text{real}$
assumes $f\ 0 = 0\ g\ 0 = 0$
assumes *summable*: $\text{summable } (\lambda n. \text{norm } (f\ n / \text{real } n\ \text{powr } s))$
 $\text{summable } (\lambda n. \text{norm } (g\ n / \text{real } n\ \text{powr } s))$
shows *summable-dirichlet-prod-real*: $\text{summable } (\lambda n. \text{norm } (\text{dirichlet-prod } f\ g\ n / \text{real } n\ \text{powr } s))$
and *suminf-dirichlet-prod-real*:
 $(\sum n. \text{dirichlet-prod } f\ g\ n / \text{real } n\ \text{powr } s) =$
 $(\sum n. f\ n / \text{real } n\ \text{powr } s) * (\sum n. g\ n / \text{real } n\ \text{powr } s)$
<proof>

lemma *fds-abs-converges-mult*:
fixes $s :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
assumes *fds-abs-converges* $f\ s\ \text{fds-abs-converges } g\ s$
shows *fds-abs-converges* $(f * g)\ s$
<proof>

lemma *fds-abs-converges-power*:
fixes $s :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
shows *fds-abs-converges* $f\ s \implies \text{fds-abs-converges } (f \wedge n)\ s$
<proof>

lemma *fds-abs-converges-prod*:
fixes $s :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
shows $(\bigwedge x. x \in A \implies \text{fds-abs-converges } (f\ x)\ s) \implies \text{fds-abs-converges } (\text{prod } f\ A)\ s$
<proof>

lemma *abs-conv-abscissa-mult-le*:
 $\text{abs-conv-abscissa } (f * g) :: 'a :: \{\text{dirichlet-series fds}\} \leq$
 $\max (\text{abs-conv-abscissa } f) (\text{abs-conv-abscissa } g)$
<proof>

lemma *abs-conv-abscissa-mult-leI*:
 $\text{abs-conv-abscissa } (f :: 'a :: \{\text{dirichlet-series fds}\}) \leq d \implies$

$abs\text{-conv-}abscissa\ g \leq d \implies abs\text{-conv-}abscissa\ (f * g) \leq d$
 ⟨proof⟩

lemma *abs-conv-abscissa-shift* [simp]:

$abs\text{-conv-}abscissa\ (fds\text{-shift}\ c\ f) = abs\text{-conv-}abscissa\ (f :: 'a :: dirichlet\text{-series}\ fds)$
 $+ c \cdot 1$
 ⟨proof⟩

lemma *eval-fds-mult*:

fixes $s :: 'a :: \{nat\text{-power}, real\text{-normed-field}, banach, second\text{-countable-topology}\}$
assumes $fds\text{-abs-converges}\ f\ s\ fds\text{-abs-converges}\ g\ s$
shows $eval\text{-fds}\ (f * g)\ s = eval\text{-fds}\ f\ s * eval\text{-fds}\ g\ s$
 ⟨proof⟩

lemma *eval-fds-power*:

fixes $s :: 'a :: \{nat\text{-power}, real\text{-normed-field}, banach, second\text{-countable-topology}\}$
assumes $fds\text{-abs-converges}\ f\ s$
shows $eval\text{-fds}\ (f \wedge^n)\ s = eval\text{-fds}\ f\ s \wedge^n$
 ⟨proof⟩

lemma *eval-fds-prod*:

fixes $s :: 'a :: \{nat\text{-power}, real\text{-normed-field}, banach, second\text{-countable-topology}\}$
assumes $(\bigwedge x. x \in A \implies fds\text{-abs-converges}\ (f\ x)\ s)$
shows $eval\text{-fds}\ (prod\ f\ A)\ s = (\prod_{x \in A}. eval\text{-fds}\ (f\ x)\ s)$ ⟨proof⟩

lemma *eval-fds-inverse*:

fixes $s :: 'a :: \{nat\text{-power}, real\text{-normed-field}, banach, second\text{-countable-topology}\}$
assumes $fds\text{-abs-converges}\ f\ s\ fds\text{-abs-converges}\ (inverse\ f)\ s\ fds\text{-nth}\ f\ 1 \neq 0$
shows $eval\text{-fds}\ (inverse\ f)\ s = inverse\ (eval\text{-fds}\ f\ s)$
 ⟨proof⟩

lemma *eval-fds-integral-has-field-derivative*:

fixes $s :: 'a :: dirichlet\text{-series}$
assumes $ereal\ (s \cdot 1) > abs\text{-conv-}abscissa\ f$
assumes $fds\text{-nth}\ f\ 1 = 0$
shows $(eval\text{-fds}\ (fds\text{-integral}\ c\ f))\ has\text{-field-derivative}\ eval\text{-fds}\ f\ s)\ (at\ s)$
 ⟨proof⟩

lemma *holomorphic-fds-eval* [holomorphic-intros]:

$A \subseteq \{z. Re\ z > conv\text{-}abscissa\ f\} \implies eval\text{-fds}\ f\ holomorphic\text{-on}\ A$
 ⟨proof⟩

lemma *analytic-fds-eval* [holomorphic-intros]:

assumes $A \subseteq \{z. Re\ z > conv\text{-}abscissa\ f\}$
shows $eval\text{-fds}\ f\ analytic\text{-on}\ A$
 ⟨proof⟩

lemma *conv-abscissa-0* [simp]:

$conv\text{-}abscissa\ (0 :: 'a :: dirichlet\text{-series}\ fds) = -\infty$

<proof>

lemma *abs-conv-abscissa-0* [*simp*]:

abs-conv-abscissa (0 :: 'a :: *dirichlet-series fds*) = $-\infty$

<proof>

lemma *conv-abscissa-1* [*simp*]:

conv-abscissa (1 :: 'a :: *dirichlet-series fds*) = $-\infty$

<proof>

lemma *abs-conv-abscissa-1* [*simp*]:

abs-conv-abscissa (1 :: 'a :: *dirichlet-series fds*) = $-\infty$

<proof>

lemma *conv-abscissa-const* [*simp*]:

conv-abscissa (*fds-const* (c :: 'a :: *dirichlet-series*)) = $-\infty$

<proof>

lemma *abs-conv-abscissa-const* [*simp*]:

abs-conv-abscissa (*fds-const* (c :: 'a :: *dirichlet-series*)) = $-\infty$

<proof>

lemma *conv-abscissa-numeral* [*simp*]:

conv-abscissa (*numeral* n :: 'a :: *dirichlet-series fds*) = $-\infty$

<proof>

lemma *abs-conv-abscissa-numeral* [*simp*]:

abs-conv-abscissa (*numeral* n :: 'a :: *dirichlet-series fds*) = $-\infty$

<proof>

lemma *abs-conv-abscissa-power-le*:

abs-conv-abscissa ($f \wedge n$:: 'a :: *dirichlet-series fds*) \leq *abs-conv-abscissa* f

<proof>

lemma *abs-conv-abscissa-power-leI*:

abs-conv-abscissa (f :: 'a :: *dirichlet-series fds*) $\leq d \implies$ *abs-conv-abscissa* ($f \wedge n$) $\leq d$

<proof>

lemma *abs-conv-abscissa-prod-le*:

assumes $\bigwedge x. x \in A \implies$ *abs-conv-abscissa* (f x :: 'a :: *dirichlet-series fds*) $\leq d$

shows *abs-conv-abscissa* (*prod* f A) $\leq d$ *<proof>*

lemma *conv-abscissa-add-le*:

conv-abscissa (f + g :: 'a :: *dirichlet-series fds*) \leq *max* (*conv-abscissa* f) (*conv-abscissa* g)

<proof>

lemma *conv-abscissa-add-leI*:

$conv-abs\ cissa (f :: 'a :: dirichlet-series\ fds) \leq d \implies conv-abs\ cissa\ g \leq d \implies$
 $conv-abs\ cissa (f + g) \leq d$
 <proof>

lemma *conv-abs\ cissa-sum-leI*:

assumes $\bigwedge x. x \in A \implies conv-abs\ cissa (f\ x :: 'a :: dirichlet-series\ fds) \leq d$
shows $conv-abs\ cissa (sum\ f\ A) \leq d$ <proof>

lemma *abs-conv-abs\ cissa-add-le*:

$abs-conv-abs\ cissa (f + g :: 'a :: dirichlet-series\ fds) \leq max (abs-conv-abs\ cissa\ f)$
 $(abs-conv-abs\ cissa\ g)$
 <proof>

lemma *abs-conv-abs\ cissa-add-leI*:

$abs-conv-abs\ cissa (f :: 'a :: dirichlet-series\ fds) \leq d \implies abs-conv-abs\ cissa\ g \leq d$
 \implies
 $abs-conv-abs\ cissa (f + g) \leq d$
 <proof>

lemma *abs-conv-abs\ cissa-sum-leI*:

assumes $\bigwedge x. x \in A \implies abs-conv-abs\ cissa (f\ x :: 'a :: dirichlet-series\ fds) \leq d$
shows $abs-conv-abs\ cissa (sum\ f\ A) \leq d$ <proof>

lemma *fds-converges-cmult-left [intro]*:

assumes *fds-converges* $f\ s$
shows *fds-converges* $(fds-const\ c * f)\ s$
 <proof>

lemma *fds-converges-cmult-right [intro]*:

assumes *fds-converges* $f\ s$
shows *fds-converges* $(f * fds-const\ c)\ s$
 <proof>

lemma *conv-abs\ cissa-cmult-left [simp]*:

fixes $c :: 'a :: dirichlet-series$ **assumes** $c \neq 0$
shows $conv-abs\ cissa (fds-const\ c * f) = conv-abs\ cissa\ f$
 <proof>

lemma *conv-abs\ cissa-cmult-right [simp]*:

fixes $c :: 'a :: dirichlet-series$ **assumes** $c \neq 0$
shows $conv-abs\ cissa (f * fds-const\ c) = conv-abs\ cissa\ f$
 <proof>

lemma *abs-conv-abs\ cissa-cmult*:

fixes $c :: 'a :: dirichlet-series$ **assumes** $c \neq 0$
shows $abs-conv-abs\ cissa (fds-const\ c * f) = abs-conv-abs\ cissa\ f$
 <proof>

lemma *conv-abs\ cissa-minus [simp]*:

fixes $f :: 'a :: \text{dirichlet-series fds}$
shows $\text{conv-absclissa } (-f) = \text{conv-absclissa } f$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-absclissa-minus}$ [simp]:
fixes $f :: 'a :: \text{dirichlet-series fds}$
shows $\text{abs-conv-absclissa } (-f) = \text{abs-conv-absclissa } f$
 $\langle \text{proof} \rangle$

lemma $\text{conv-absclissa-diff-le}$:
 $\text{conv-absclissa } (f - g :: 'a :: \text{dirichlet-series fds}) \leq \max (\text{conv-absclissa } f) (\text{conv-absclissa } g)$
 $\langle \text{proof} \rangle$

lemma $\text{conv-absclissa-diff-leI}$:
 $\text{conv-absclissa } (f :: 'a :: \text{dirichlet-series fds}) \leq d \implies \text{conv-absclissa } g \leq d \implies$
 $\text{conv-absclissa } (f - g) \leq d$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-absclissa-diff-le}$:
 $\text{abs-conv-absclissa } (f - g :: 'a :: \text{dirichlet-series fds}) \leq$
 $\max (\text{abs-conv-absclissa } f) (\text{abs-conv-absclissa } g)$
 $\langle \text{proof} \rangle$

lemma $\text{abs-conv-absclissa-diff-leI}$:
 $\text{abs-conv-absclissa } (f :: 'a :: \text{dirichlet-series fds}) \leq d \implies \text{abs-conv-absclissa } g \leq d$
 \implies
 $\text{abs-conv-absclissa } (f - g) \leq d$
 $\langle \text{proof} \rangle$

lemmas $\text{eval-fds-integral-has-field-derivative}'$ [derivative-intros] =
 $\text{DERIV-chain}'[\text{OF } - \text{eval-fds-integral-has-field-derivative}]$

lemma $\text{abs-conv-absclissa-completely-multiplicative-log-deriv}$:
fixes $f :: 'a :: \text{dirichlet-series fds}$
assumes $\text{completely-multiplicative-function } (\text{fds-nth } f) \text{ fds-nth } f \ 1 \neq 0$
shows $\text{abs-conv-absclissa } (\text{fds-deriv } f / f) \leq \text{abs-conv-absclissa } f$
 $\langle \text{proof} \rangle$

12.4 Uniqueness

context
assumes $\text{SORT-CONSTRAINT } ('a :: \text{dirichlet-series})$
begin

lemma $\text{norm-dirichlet-series-cutoff-le}$:
assumes $\text{fds-abs-converges } f (s0 :: 'a) \ N > 0 \ s \cdot 1 \geq c \ c \geq s0 \cdot 1$
shows $\text{summable } (\lambda n. \text{fds-nth } f (n + N) / \text{nat-power } (n + N) \ s)$
 $\text{summable } (\lambda n. \text{norm } (\text{fds-nth } f (n + N)) / \text{nat-power } (n + N) \ c)$

and $norm (\sum n. fds\text{-}nth\ f\ (n + N) / nat\text{-}power\ (n + N)\ s) \leq$
 $(\sum n. norm\ (fds\text{-}nth\ f\ (n + N)) / nat\text{-}power\ (n + N)\ c) / nat\text{-}power$
 $N\ (s \cdot 1 - c)$
 $\langle proof \rangle$

lemma *eval-fds-zeroD-aux*:
fixes $h :: 'a\ fds$
assumes *conv*: $fds\text{-}abs\text{-}converges\ h\ (s0 :: 'a)$
assumes *freq*: $frequently\ (\lambda s. eval\text{-}fds\ h\ s = 0)\ ((\lambda s. s \cdot 1)\ going\text{-}to\ at\text{-}top)$
shows $h = 0$
 $\langle proof \rangle$

lemma *eval-fds-zeroD*:
fixes $h :: 'a\ fds$
assumes *conv*: $conv\text{-}abscissa\ h < \infty$
assumes *freq*: $frequently\ (\lambda s. eval\text{-}fds\ h\ s = 0)\ ((\lambda s. s \cdot 1)\ going\text{-}to\ at\text{-}top)$
shows $h = 0$
 $\langle proof \rangle$

lemma *eval-fds-eqD*:
fixes $f\ g :: 'a\ fds$
assumes *conv*: $conv\text{-}abscissa\ f < \infty\ conv\text{-}abscissa\ g < \infty$
assumes *eq*: $frequently\ (\lambda s. eval\text{-}fds\ f\ s = eval\text{-}fds\ g\ s)\ ((\lambda s. s \cdot 1)\ going\text{-}to\ at\text{-}top)$
shows $f = g$
 $\langle proof \rangle$

end

12.5 Limit at infinity

lemma *eval-fds-at-top-tail-bound*:
fixes $f :: 'a :: dirichlet\text{-}series\ fds$
assumes *c*: $ereal\ c > abs\text{-}conv\text{-}abscissa\ f$
defines $B \equiv (\sum n. norm\ (fds\text{-}nth\ f\ (n+2)) / real\ (n+2)\ powr\ c) * 2\ powr\ c$
assumes *s*: $s \cdot 1 \geq c$
shows $norm\ (eval\text{-}fds\ f\ s - fds\text{-}nth\ f\ 1) \leq B / 2\ powr\ (s \cdot 1)$
 $\langle proof \rangle$

lemma *tendsto-eval-fds-Re-at-top*:
assumes *conv-abscissa* ($f :: 'a :: dirichlet\text{-}series\ fds$) $\neq \infty$
assumes *lim*: $filterlim\ (\lambda x. S\ x \cdot 1)\ at\text{-}top\ F$
shows $((\lambda x. eval\text{-}fds\ f\ (S\ x)) \longrightarrow fds\text{-}nth\ f\ 1)\ F$
 $\langle proof \rangle$

lemma *tendsto-eval-fds-Re-at-top'*:
assumes *conv-abscissa* ($f :: complex\ fds$) $\neq \infty$
shows $uniform\text{-}limit\ UNIV\ (\lambda\sigma\ t. eval\text{-}fds\ f\ (of\text{-}real\ \sigma + of\text{-}real\ t * i))$
 $(\lambda\cdot. fds\text{-}nth\ f\ 1)\ at\text{-}top$

<proof>

lemma *tendsto-eval-fds-Re-going-to-at-top*:

assumes *conv-abscissa* ($f :: 'a :: \text{dirichlet-series fds}$) $\neq \infty$

shows $((\lambda s. \text{eval-fds } f \ s) \longrightarrow \text{fds-nth } f \ 1) ((\lambda s. s \cdot 1) \text{ going-to at-top})$

<proof>

lemma *tendsto-eval-fds-Re-going-to-at-top'*:

assumes *conv-abscissa* ($f :: \text{complex fds}$) $\neq \infty$

shows $((\lambda s. \text{eval-fds } f \ s) \longrightarrow \text{fds-nth } f \ 1) (\text{Re going-to at-top})$

<proof>

Any Dirichlet series that is not identically zero and does not diverge everywhere has a half-plane in which it converges and is non-zero.

theorem *fds-nonzero-halfplane-exists*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

assumes *conv-abscissa* $f < \infty$ $f \neq 0$

shows *eventually* $(\lambda s. \text{fds-converges } f \ s \wedge \text{eval-fds } f \ s \neq 0) ((\lambda s. s \cdot 1) \text{ going-to at-top})$

<proof>

12.6 Normed series

lemma *fds-converges-norm-iff [simp]*:

fixes $s :: 'a :: \{\text{nat-power-normed-field, banach}\}$

shows $\text{fds-converges } (\text{fds-norm } f) \ (s \cdot 1) \longleftrightarrow \text{fds-abs-converges } f \ s$

<proof>

lemma *fds-abs-converges-norm-iff [simp]*:

fixes $s :: 'a :: \{\text{nat-power-normed-field, banach}\}$

shows $\text{fds-abs-converges } (\text{fds-norm } f) \ (s \cdot 1) \longleftrightarrow \text{fds-abs-converges } f \ s$

<proof>

lemma *fds-converges-norm-iff'*:

fixes $f :: 'a :: \{\text{nat-power-normed-field, banach}\} \text{ fds}$

shows $\text{fds-converges } (\text{fds-norm } f) \ s \longleftrightarrow \text{fds-abs-converges } f \ (\text{of-real } s)$

<proof>

lemma *fds-abs-converges-norm-iff'*:

fixes $f :: 'a :: \{\text{nat-power-normed-field, banach}\} \text{ fds}$

shows $\text{fds-abs-converges } (\text{fds-norm } f) \ s \longleftrightarrow \text{fds-abs-converges } f \ (\text{of-real } s)$

<proof>

lemma *abs-conv-abscissa-norm [simp]*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

shows $\text{abs-conv-abscissa } (\text{fds-norm } f) = \text{abs-conv-abscissa } f$

<proof>

lemma *conv-abscissa-norm [simp]*:

fixes $f :: 'a :: \text{dirichlet-series fds}$
shows $\text{conv-abscissa (fds-norm } f) = \text{abs-conv-abscissa } f$
 $\langle \text{proof} \rangle$

lemma

fixes $f g :: 'a :: \text{dirichlet-series fds}$
assumes $\text{fds-abs-converges (fds-norm } f) s \text{ fds-abs-converges (fds-norm } g) s$
shows $\text{fds-abs-converges-norm-mult: fds-abs-converges (fds-norm (} f * g)) s$
and $\text{eval-fds-norm-mult-le:}$
 $\text{eval-fds (fds-norm (} f * g)) s \leq \text{eval-fds (fds-norm } f) s * \text{eval-fds (fds-norm}$
 $g) s$
 $\langle \text{proof} \rangle$

lemma $\text{eval-fds-norm-nonneg:}$

assumes $\text{fds-abs-converges (fds-norm } f) s$
shows $\text{eval-fds (fds-norm } f) s \geq 0$
 $\langle \text{proof} \rangle$

lemma

fixes $f :: 'a :: \text{dirichlet-series fds}$
assumes $\text{fds-abs-converges (fds-norm } f) s$
shows $\text{fds-abs-converges-norm-power: fds-abs-converges (fds-norm (} f \wedge n)) s$
and $\text{eval-fds-norm-power-le:}$
 $\text{eval-fds (fds-norm (} f \wedge n)) s \leq \text{eval-fds (fds-norm } f) s \wedge n$
 $\langle \text{proof} \rangle$

12.7 Logarithms of Dirichlet series

lemma $\text{eventually-gt-ereal-at-top: } c \neq \infty \implies \text{eventually } (\lambda x. \text{ereal } x > c) \text{ at-top}$
 $\langle \text{proof} \rangle$

lemma $\text{eval-fds-log-deriv:}$

fixes $s :: 'a :: \text{dirichlet-series}$
assumes $\text{fds-nth } f \ 1 \neq 0 \ s \cdot 1 > \text{abs-conv-abscissa } f$
 $s \cdot 1 > \text{abs-conv-abscissa (fds-deriv } f / f)$
assumes $\text{eval-fds } f \ s \neq 0$
shows $\text{eval-fds (fds-deriv } f / f) s = \text{eval-fds (fds-deriv } f) s / \text{eval-fds } f \ s$
 $\langle \text{proof} \rangle$

Given a sufficiently nice absolutely convergent Dirichlet series that converges to some function $f(s)$ and a holomorphic branch of $\ln f(s)$, we can construct a Dirichlet series that absolutely converges to that logarithm.

lemma eval-fds-ln:

fixes $s0 :: \text{ereal}$
assumes $\text{nz: } \bigwedge s. \text{Re } s > s0 \implies \text{eval-fds } f \ s \neq 0 \ \text{fds-nth } f \ 1 \neq 0$
assumes $l: \text{exp } l = \text{fds-nth } f \ 1 \ ((g \circ \text{of-real}) \longrightarrow l) \text{ at-top}$
assumes $g: \bigwedge s. \text{Re } s > s0 \implies \text{exp } (g \ s) = \text{eval-fds } f \ s$
assumes $\text{holo-g: } g \text{ holomorphic-on } \{s. \text{Re } s > s0\}$
assumes $\text{ereal } (\text{Re } s) > s0$

assumes $s0 \geq \text{abs-conv-abscissa } f$ **and** $s0 \geq \text{abs-conv-abscissa } (f\text{-deriv } f / f)$
shows $\text{eval-fds } (f\text{-ln } l f) s = g s$
 $\langle \text{proof} \rangle$

Less explicitly: For a sufficiently nice absolutely convergent Dirichlet series converging to a function $f(s)$, the formal logarithm absolutely converges to some logarithm of $f(s)$.

lemma *eval-fds-ln'*:

fixes $s0 :: \text{ereal}$
assumes $\text{ereal } (\text{Re } s) > s0$
assumes $s0 \geq \text{abs-conv-abscissa } f$ **and** $s0 \geq \text{abs-conv-abscissa } (f\text{-deriv } f / f)$
and $\text{nz}: \bigwedge s. \text{Re } s > s0 \implies \text{eval-fds } f s \neq 0 \text{ fds-nth } f 1 \neq 0$
assumes $l: \text{exp } l = \text{fds-nth } f 1$
shows $\text{exp } (\text{eval-fds } (f\text{-ln } l f) s) = \text{eval-fds } f s$
 $\langle \text{proof} \rangle$

lemma *fds-ln-completely-multiplicative*:

fixes $f :: 'a :: \text{dirichlet-series fds}$
assumes *completely-multiplicative-function* $(\text{fds-nth } f)$
assumes $\text{fds-nth } f 1 \neq 0$
shows $\text{fds-ln } l f = \text{fds } (\lambda n. \text{if } n = 1 \text{ then } l \text{ else } \text{fds-nth } f n * \text{mangoldt } n /_R \text{ln } n)$
 $\langle \text{proof} \rangle$

lemma *eval-fds-ln-completely-multiplicative-strong*:

fixes $s :: 'a :: \text{dirichlet-series}$ **and** $l :: 'a$ **and** $f :: 'a \text{ fds}$ **and** $g :: \text{nat} \Rightarrow 'a$
defines $h \equiv \text{fds } (\lambda n. \text{fds-nth } (f\text{-ln } l f) n * g n)$
assumes *fds-abs-converges* $h s$
assumes *completely-multiplicative-function* $(\text{fds-nth } f)$ **and** $\text{fds-nth } f 1 \neq 0$
shows $(\lambda(p,k). (\text{fds-nth } f p / \text{nat-power } p s) \wedge \text{Suc } k * g (p \wedge \text{Suc } k) / \text{of-nat } (\text{Suc } k))$
 $\text{abs-summable-on } (\{p. \text{prime } p\} \times \text{UNIV})$ **(is ?th1)**
and $\text{eval-fds } h s = l * g 1 + (\sum_a(p, k) \in \{p. \text{prime } p\} \times \text{UNIV}. (\text{fds-nth } f p / \text{nat-power } p s) \wedge \text{Suc } k * g (p \wedge \text{Suc } k) / \text{of-nat } (\text{Suc } k))$
(is ?th2)
 $\langle \text{proof} \rangle$

lemma *eval-fds-ln-completely-multiplicative*:

fixes $s :: 'a :: \text{dirichlet-series}$ **and** $l :: 'a$ **and** $f :: 'a \text{ fds}$
assumes *completely-multiplicative-function* $(\text{fds-nth } f)$ **and** $\text{fds-nth } f 1 \neq 0$
assumes $s \cdot 1 > \text{abs-conv-abscissa } (f\text{-deriv } f / f)$
shows $(\lambda(p,k). (\text{fds-nth } f p / \text{nat-power } p s) \wedge \text{Suc } k / \text{of-nat } (\text{Suc } k))$
 $\text{abs-summable-on } (\{p. \text{prime } p\} \times \text{UNIV})$ **(is ?th1)**
and $\text{eval-fds } (f\text{-ln } l f) s =$
 $l + (\sum_a(p, k) \in \{p. \text{prime } p\} \times \text{UNIV}. (\text{fds-nth } f p / \text{nat-power } p s) \wedge \text{Suc } k / \text{of-nat } (\text{Suc } k))$ **(is ?th2)**
 $\langle \text{proof} \rangle$

12.8 Exponential and logarithm

lemma *summable-fds-exp-aux*:

assumes $\text{fds-nth } f' \ 1 = (0 :: 'a :: \text{real-normed-algebra-1})$

shows $\text{summable } (\lambda k. \text{fds-nth } (f' \ ^k) \ n \ /_R \ \text{fact } k)$

<proof>

lemma

fixes $f :: 'a :: \text{dirichlet-series fds}$

assumes $\text{fds-abs-converges } f \ s$

shows $\text{fds-abs-converges-exp: fds-abs-converges } (\text{fds-exp } f) \ s$

and $\text{eval-fds-exp: eval-fds } (\text{fds-exp } f) \ s = \text{exp } (\text{eval-fds } f \ s)$

<proof>

lemma *fds-exp-add*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

shows $\text{fds-exp } (f + g) = \text{fds-exp } f * \text{fds-exp } g$

<proof>

lemma *fds-exp-minus*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

shows $\text{fds-exp } (-f) = \text{inverse } (\text{fds-exp } f)$

<proof>

lemma *abs-conv-abscissa-exp*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

shows $\text{abs-conv-abscissa } (\text{fds-exp } f) \leq \text{abs-conv-abscissa } f$

<proof>

lemma *fds-deriv-exp [simp]*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

shows $\text{fds-deriv } (\text{fds-exp } f) = \text{fds-exp } f * \text{fds-deriv } f$

<proof>

lemma *fds-exp-ln-strong*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

assumes $\text{fds-nth } f \ (\text{Suc } 0) \neq 0$

shows $\text{fds-exp } (\text{fds-ln } l \ f) = \text{fds-const } (\text{exp } l \ / \ \text{fds-nth } f \ (\text{Suc } 0)) * f$

<proof>

lemma *fds-exp-ln [simp]*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

assumes $\text{exp } l = \text{fds-nth } f \ (\text{Suc } 0)$

shows $\text{fds-exp } (\text{fds-ln } l \ f) = f$

<proof>

lemma *fds-ln-exp [simp]*:

fixes $f :: 'a :: \text{dirichlet-series fds}$

assumes $l = \text{fds-nth } f \ (\text{Suc } 0)$

shows $\text{fds-ln } l \ (\text{fds-exp } f) = f$

<proof>

12.9 Euler products

lemma *fds-euler-product-LIMSEQ*:

fixes $f :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
fds

assumes *multiplicative-function* ($\text{fds-nth } f$) **and** *fds-abs-converges* $f\ s$

shows $(\lambda n. \prod_{p \leq n}. \text{if prime } p \text{ then } \sum i. \text{fds-nth } f\ (p \wedge i) / \text{nat-power } (p \wedge i) \text{ else } 1) \longrightarrow$

eval-fds $f\ s$

<proof>

lemma *fds-euler-product-LIMSEQ'*:

fixes $f :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
fds

assumes *completely-multiplicative-function* ($\text{fds-nth } f$) **and** *fds-abs-converges* $f\ s$

shows $(\lambda n. \prod_{p \leq n}. \text{if prime } p \text{ then inverse } (1 - \text{fds-nth } f\ p / \text{nat-power } p\ s) \text{ else } 1) \longrightarrow$

eval-fds $f\ s$

<proof>

lemma *fds-abs-convergent-euler-product*:

fixes $f :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
fds

assumes *multiplicative-function* ($\text{fds-nth } f$) **and** *fds-abs-converges* $f\ s$

shows *abs-convergent-prod*

$(\lambda p. \text{if prime } p \text{ then } \sum i. \text{fds-nth } f\ (p \wedge i) / \text{nat-power } (p \wedge i) \text{ else } 1)$

<proof>

lemma *fds-abs-convergent-euler-product'*:

fixes $f :: 'a :: \{\text{nat-power, real-normed-field, banach, second-countable-topology}\}$
fds

assumes *completely-multiplicative-function* ($\text{fds-nth } f$) **and** *fds-abs-converges* $f\ s$

shows *abs-convergent-prod*

$(\lambda p. \text{if prime } p \text{ then inverse } (1 - \text{fds-nth } f\ p / \text{nat-power } p\ s) \text{ else } 1)$

<proof>

lemma *fds-abs-convergent-zero-iff*:

fixes $f :: 'a :: \{\text{nat-power-field, real-normed-field, banach, second-countable-topology}\}$
fds

assumes *completely-multiplicative-function* ($\text{fds-nth } f$)

assumes *fds-abs-converges* $f\ s$

shows $\text{eval-fds } f\ s = 0 \iff (\exists p. \text{prime } p \wedge \text{fds-nth } f\ p = \text{nat-power } p\ s)$

<proof>

lemma

fixes $s :: 'a :: \{\text{nat-power-normed-field, banach, euclidean-space}\}$

assumes $s \cdot 1 > 1$

shows *euler-product-fds-zeta*:
 $(\lambda n. \prod_{p \leq n} \text{if prime } p \text{ then inverse } (1 - 1 / \text{nat-power } p \ s) \text{ else } 1)$
 $\longrightarrow \text{eval-fds fds-zeta } s \text{ (is ?th1)}$
and *eval-fds-zeta-nonzero*: $\text{eval-fds fds-zeta } s \neq 0$
 $\langle \text{proof} \rangle$

lemma *fds-primepow-subseries-euler-product-cm*:
fixes $f :: 'a :: \text{dirichlet-series fds}$
assumes *completely-multiplicative-function* $(\text{fds-nth } f)$ *prime* p
assumes $s \cdot 1 > \text{abs-conv-abscissa } f$
shows $\text{eval-fds } (\text{fds-primepow-subseries } p \ f) \ s = 1 / (1 - \text{fds-nth } f \ p / \text{nat-power } p \ s)$
 $\langle \text{proof} \rangle$

12.10 Non-negative Dirichlet series

lemma *nonneg-Reals-sum*: $(\bigwedge x. x \in A \implies f \ x \in \mathbb{R}_{\geq 0}) \implies \text{sum } f \ A \in \mathbb{R}_{\geq 0}$
 $\langle \text{proof} \rangle$

locale *nonneg-dirichlet-series* =
fixes $f :: 'a :: \text{dirichlet-series fds}$
assumes *nonneg-coeffs-aux*: $n > 0 \implies \text{fds-nth } f \ n \in \mathbb{R}_{\geq 0}$
begin

lemma *nonneg-coeffs*: $\text{fds-nth } f \ n \in \mathbb{R}_{\geq 0}$
 $\langle \text{proof} \rangle$

end

lemma *nonneg-dirichlet-series-0* [*simp,intro*]: *nonneg-dirichlet-series 0*
 $\langle \text{proof} \rangle$

lemma *nonneg-dirichlet-series-1* [*simp,intro*]: *nonneg-dirichlet-series 1*
 $\langle \text{proof} \rangle$

lemma *nonneg-dirichlet-series-const* [*simp,intro*]:
 $c \in \mathbb{R}_{\geq 0} \implies \text{nonneg-dirichlet-series } (\text{fds-const } c)$
 $\langle \text{proof} \rangle$

lemma *nonneg-dirichlet-series-add* [*intro*]:
assumes *nonneg-dirichlet-series f nonneg-dirichlet-series g*
shows *nonneg-dirichlet-series (f + g)*
 $\langle \text{proof} \rangle$

lemma *nonneg-dirichlet-series-mult* [*intro*]:
assumes *nonneg-dirichlet-series f nonneg-dirichlet-series g*
shows *nonneg-dirichlet-series (f * g)*
 $\langle \text{proof} \rangle$

lemma *nonneg-dirichlet-series-power* [intro]:
assumes *nonneg-dirichlet-series* f
shows *nonneg-dirichlet-series* $(f \wedge n)$
 ⟨proof⟩

context *nonneg-dirichlet-series*
begin

lemma *nonneg-exp* [intro]: *nonneg-dirichlet-series* (*fds-exp* f)
 ⟨proof⟩

end

lemma *nonneg-dirichlet-series-lnD*:
assumes *nonneg-dirichlet-series* (*fds-ln* l f) *exp* $l = \text{fds-nth } f \text{ (Suc } 0)$
shows *nonneg-dirichlet-series* f
 ⟨proof⟩

context *nonneg-dirichlet-series*
begin

lemma *fds-of-real-norm*: *fds-of-real* (*fds-norm* f) = f
 ⟨proof⟩

end

lemma *pringsheim-landau-aux*:
fixes $c :: \text{real}$ **and** $f :: \text{complex fds}$
assumes *nonneg-dirichlet-series* f
assumes *abscissa*: $c \geq \text{abs-conv-abscissa } f$
assumes $g: \bigwedge s. s \in A \implies \text{Re } s > c \implies g \ s = \text{eval-fds } f \ s$
assumes g *holomorphic-on* A *open* A $c \in A$
shows $\exists x. x < c \wedge \text{fds-abs-converges } f \text{ (of-real } x)$
 ⟨proof⟩

theorem *pringsheim-landau*:
fixes $c :: \text{real}$ **and** $f :: \text{complex fds}$
assumes *nonneg-dirichlet-series* f
assumes *abscissa*: $\text{abs-conv-abscissa } f = c$
assumes $g: \bigwedge s. s \in A \implies \text{Re } s > c \implies g \ s = \text{eval-fds } f \ s$
assumes g *holomorphic-on* A *open* A $c \in A$
shows *False*
 ⟨proof⟩

corollary *entire-continuation-imp-abs-conv-abscissa-MInfty*:
assumes *nonneg-dirichlet-series* f
assumes $c: c \geq \text{abs-conv-abscissa } f$
assumes $g: \bigwedge s. \text{Re } s > c \implies g \ s = \text{eval-fds } f \ s$

assumes *holo*: g holomorphic-on UNIV
shows *abs-conv-abscissa* $f = -\infty$
 ⟨*proof*⟩

12.11 Convergence of the ζ and Möbius μ series

lemma *fds-abs-summable-zeta-real-iff* [*simp*]:
fds-abs-converges fds-zeta $s \longleftrightarrow s > (1 :: \text{real})$
 ⟨*proof*⟩

lemma *fds-abs-summable-zeta-real*: $s > (1 :: \text{real}) \implies \text{fds-abs-converges fds-zeta}$
 s
 ⟨*proof*⟩

lemma *fds-abs-converges-moebius-mu-real*:
assumes $s > (1 :: \text{real})$
shows *fds-abs-converges (fds moebius-mu)* s
 ⟨*proof*⟩

12.12 Application to the Möbius μ function

lemma *inverse-squares-sums'*: $(\lambda n. 1 / \text{real } n^2) \text{ sums } (\pi^2 / 6)$
 ⟨*proof*⟩

lemma *norm-summable-moebius-over-square*:
summable $(\lambda n. \text{norm } (\text{moebius-mu } n / \text{real } n^2))$
 ⟨*proof*⟩

lemma *summable-moebius-over-square*:
summable $(\lambda n. \text{moebius-mu } n / \text{real } n^2)$
 ⟨*proof*⟩

lemma *moebius-over-square-sums*: $(\lambda n. \text{moebius-mu } n / n^2) \text{ sums } (6 / \pi^2)$
 ⟨*proof*⟩

end

13 Asymptotics of summatory arithmetic functions

theory *Arithmetic-Summatory-Asymptotics*
imports
Euler-MacLaurin.Euler-MacLaurin-Landau
Arithmetic-Summatory
Dirichlet-Series-Analysis
Landau-Symbols.Landau-More
begin

13.1 Auxiliary bounds

lemma *sum-inverse-squares-tail-bound*:

assumes $d > 0$

shows $\text{summable } (\lambda n. 1 / (\text{real } (\text{Suc } n) + d) ^ 2)$
 $(\sum n. 1 / (\text{real } (\text{Suc } n) + d) ^ 2) \leq 1 / d$

<proof>

lemma *moebius-sum-tail-bound*:

assumes $d > 0$

shows $\text{abs } (\sum n. \text{moebius-mu } (\text{Suc } n + d) / \text{real } (\text{Suc } n + d) ^ 2) \leq 1 / d$ (**is**
 $\text{abs } ?S \leq -$)

<proof>

lemma *sum-upto-inverse-bound*:

$\text{sum-upto } (\lambda i. 1 / \text{real } i) x \geq 0$

$\text{eventually } (\lambda x. \text{sum-upto } (\lambda i. 1 / \text{real } i) x \leq \ln x + 13 / 22)$ *at-top*

<proof>

lemma *sum-upto-inverse-bigo*: $\text{sum-upto } (\lambda i. 1 / \text{real } i) \in O(\lambda x. \ln x)$

<proof>

lemma

defines $G \equiv (\lambda x::\text{real}. (\sum n. \text{moebius-mu } (n + \text{Suc } (\text{nat } \lfloor x \rfloor)) / (n + \text{Suc } (\text{nat } \lfloor x \rfloor)) ^ 2) :: \text{real})$

shows *moebius-sum-tail-bound'*: $\bigwedge t. t \geq 2 \implies |G t| \leq 1 / (t - 1)$

and *moebius-sum-tail-bigo*: $G \in O(\lambda t. 1 / t)$

<proof>

13.2 Summatory totient function

theorem *summatory-totient-asymptotics*:

$(\lambda x. \text{sum-upto } (\lambda n. \text{real } (\text{totient } n)) x - 3 / \pi^2 * x^2) \in O(\lambda x. x * \ln x)$

<proof>

theorem *summatory-totient-asymptotics'*:

$(\lambda x. \text{sum-upto } (\lambda n. \text{real } (\text{totient } n)) x) = o(\lambda x. 3 / \pi^2 * x^2) + o O(\lambda x. x * \ln x)$

<proof>

theorem *summatory-totient-asymptotics''*:

$\text{sum-upto } (\lambda n. \text{real } (\text{totient } n)) \sim_{[\text{at-top}]} (\lambda x. 3 / \pi^2 * x^2)$

<proof>

13.3 Asymptotic distribution of squarefree numbers

lemma *le-sqrt-iff*: $x \geq 0 \implies x \leq \text{sqrt } y \iff x^2 \leq y$

<proof>

theorem *squarefree-asymptotics*: $(\lambda x. \text{card } \{n. \text{real } n \leq x \wedge \text{squarefree } n\} - 6 / \pi^2 * x) \in O(\text{sqrt})$

<proof>

theorem *squarefree-asymptotics'*:

$(\lambda x. \text{card } \{n. \text{real } n \leq x \wedge \text{squarefree } n\}) = o(\lambda x. 6 / \pi^2 * x) + o O(\lambda x. \text{sqrt } x)$
<proof>

theorem *squarefree-asymptotics''*:

$(\lambda x. \text{card } \{n. \text{real } n \leq x \wedge \text{squarefree } n\}) \sim[\text{at-top}] (\lambda x. 6 / \pi^2 * x)$
<proof>

13.4 The hyperbola method

lemma *hyperbola-method-bigo*:

fixes $f g :: \text{nat} \Rightarrow 'a :: \text{real-normed-field}$

assumes $(\lambda x. \text{sum-upto } (\lambda n. f n * \text{sum-upto } g (x / \text{real } n)) (\text{sqrt } x) - R x) \in O(b)$

assumes $(\lambda x. \text{sum-upto } (\lambda n. \text{sum-upto } f (x / \text{real } n) * g n) (\text{sqrt } x) - S x) \in O(b)$

assumes $(\lambda x. \text{sum-upto } f (\text{sqrt } x) * \text{sum-upto } g (\text{sqrt } x) - T x) \in O(b)$

shows $(\lambda x. \text{sum-upto } (\text{dirichlet-prod } f g) x - (R x + S x - T x)) \in O(b)$
<proof>

lemma *frac-le-1*: $\text{frac } x \leq 1$

<proof>

lemma *ln-minus-ln-floor-bound*:

assumes $x \geq 2$

shows $\ln x - \ln (\text{floor } x) \in \{0..<1 / (x - 1)\}$

<proof>

lemma *ln-minus-ln-floor-bigo*:

$(\lambda x::\text{real}. \ln x - \ln (\text{floor } x)) \in O(\lambda x. 1 / x)$

<proof>

lemma *divisor-count-asymptotics-aux*:

$(\lambda x. \text{sum-upto } (\lambda n. \text{sum-upto } (\lambda-. 1) (x / \text{real } n)) (\text{sqrt } x) - (x * \ln x / 2 + \text{euler-mascheroni} * x)) \in O(\text{sqrt})$

<proof>

lemma *sum-upto-sqrt-bound*:

assumes $x \geq (0 :: \text{real})$

shows $\text{norm } ((\text{sum-upto } (\lambda-. 1) (\text{sqrt } x))^2 - x) \leq 2 * \text{norm } (\text{sqrt } x)$

<proof>

lemma *summatory-divisor-count-asymptotics*:

$(\lambda x. \text{sum-upto } (\lambda n. \text{real } (\text{divisor-count } n)) x - (x * \ln x + (2 * \text{euler-mascheroni} - 1) * x)) \in O(\text{sqrt})$

<proof>

theorem *summatory-divisor-count-asymptotics'*:

$(\lambda x. \text{sum-upto } (\lambda n. \text{real } (\text{divisor-count } n)) x) = o$
 $(\lambda x. x * \ln x + (2 * \text{euler-mascheroni} - 1) * x) + o O(\lambda x. \text{sqrt } x)$
 $\langle \text{proof} \rangle$

theorem *summatory-divisor-count-asymptotics''*:

$\text{sum-upto } (\lambda n. \text{real } (\text{divisor-count } n)) \sim[\text{at-top}] (\lambda x. x * \ln x)$
 $\langle \text{proof} \rangle$

lemma *summatory-divisor-eq*:

$\text{sum-upto } (\lambda n. \text{real } (\text{divisor-count } n)) (\text{real } m) = \text{card } \{(n,d). n \in \{0 <..m\} \wedge d \text{ dvd } n\}$
 $\langle \text{proof} \rangle$

context

fixes $M :: \text{nat} \Rightarrow \text{real}$

defines $M \equiv \lambda m. \text{card } \{(n,d). n \in \{0 <..m\} \wedge d \text{ dvd } n\} / \text{card } \{0 <..m\}$

begin

lemma *mean-divisor-count-asymptotics*:

$(\lambda m. M m - (\ln m + 2 * \text{euler-mascheroni} - 1)) \in O(\lambda m. 1 / \text{sqrt } m)$
 $\langle \text{proof} \rangle$

theorem *mean-divisor-count-asymptotics'*:

$M = o (\lambda x. \ln x + 2 * \text{euler-mascheroni} - 1) + o O(\lambda x. 1 / \text{sqrt } x)$
 $\langle \text{proof} \rangle$

theorem *mean-divisor-count-asymptotics''*: $M \sim[\text{at-top}] \ln$

$\langle \text{proof} \rangle$

end

13.5 The asymptotic ditribution of coprime pairs

context

fixes $A :: \text{nat} \Rightarrow (\text{nat} \times \text{nat}) \text{ set}$

defines $A \equiv (\lambda N. \{(m,n) \in \{1..N\} \times \{1..N\}. \text{coprime } m n\})$

begin

lemma *coprime-pairs-asymptotics*:

$(\lambda N. \text{real } (\text{card } (A N)) - 6 / \pi^2 * (\text{real } N)^2) \in O(\lambda N. \text{real } N * \ln (\text{real } N))$
 $\langle \text{proof} \rangle$

theorem *coprime-pairs-asymptotics'*:

$(\lambda N. \text{real } (\text{card } (A N))) = o (\lambda N. 6 / \pi^2 * (\text{real } N)^2) + o O(\lambda N. \text{real } N * \ln (\text{real } N))$
 $\langle \text{proof} \rangle$

theorem *coprime-pairs-asymptotics''*:

$(\lambda N. \text{real} (\text{card} (A N))) \sim_{[at-top]} (\lambda N. 6 / \pi^2 * (\text{real } N)^2)$
 ⟨proof⟩

theorem coprime-probability-tendsto:

$(\lambda N. \text{card} (A N) / \text{card} (\{1..N\} \times \{1..N\})) \longrightarrow 6 / \pi^2$
 ⟨proof⟩

end

13.6 The asymptotics of the number of Farey fractions

definition farey-fractions :: $\text{nat} \Rightarrow \text{rat set}$ **where**

$\text{farey-fractions } N = \{q :: \text{rat} \in \{0 < .. 1\}. \text{snd} (\text{quotient-of } q) \leq \text{int } N\}$

lemma Fract-eq-coprime:

assumes $\text{Rat.Fract } a \ b = \text{Rat.Fract } c \ d \ b > 0 \ d > 0 \ \text{coprime } a \ b \ \text{coprime } c \ d$

shows $a = c \ b = d$

⟨proof⟩

lemma quotient-of-split:

$P (\text{quotient-of } q) = (\forall a \ b. \ b > 0 \longrightarrow \text{coprime } a \ b \longrightarrow q = \text{Rat.Fract } a \ b \longrightarrow P$
 $(a, b))$

⟨proof⟩

lemma quotient-of-split-asm:

$P (\text{Rat.quotient-of } q) = (\neg(\exists a \ b. \ b > 0 \wedge \text{coprime } a \ b \wedge q = \text{Rat.Fract } a \ b \wedge$
 $\neg P (a, b)))$

⟨proof⟩

lemma farey-fractions-bij:

$\text{bij-betw } (\lambda(a,b). \text{Rat.Fract } (\text{int } a) (\text{int } b))$

$\{(a,b) \mid a \ b. \ 0 < a \wedge a \leq b \wedge b \leq N \wedge \text{coprime } a \ b\} (\text{farey-fractions } N)$

⟨proof⟩

lemma card-farey-fractions: $\text{card} (\text{farey-fractions } N) = \text{sum totient } \{0 < .. N\}$

⟨proof⟩

lemma card-farey-fractions-asymptotics:

$(\lambda N. \text{real} (\text{card} (\text{farey-fractions } N)) - 3 / \pi^2 * (\text{real } N)^2) \in O(\lambda N. \text{real } N * \ln$
 $(\text{real } N))$

⟨proof⟩

theorem card-farey-fractions-asymptotics':

$(\lambda N. \text{card} (\text{farey-fractions } N)) = o (\lambda N. 3 / \pi^2 * N^2) + o O(\lambda N. N * \ln N)$

⟨proof⟩

theorem card-farey-fractions-asymptotics'':

$(\lambda N. \text{real} (\text{card} (\text{farey-fractions } N))) \sim_{[at-top]} (\lambda N. 3 / \pi^2 * (\text{real } N)^2)$

⟨proof⟩

end

14 Efficient code for number-theoretic functions

theory *Dirichlet-Efficient-Code*

imports

Main

Moebius-Mu

More-Totient

Divisor-Count

Liouville-Lambda

HOL-Library.Code-Target-Numeral

Polynomial-Factorization.Prime-Factorization

begin

definition *prime-factorization-nat'* :: *nat* \Rightarrow (*nat* \times *nat*) *list* **where**

prime-factorization-nat' n = (
 let ps = *prime-factorization-nat n*
 in map ($\lambda p. (p, \text{length } (\text{filter } ((=) p) ps) - 1)$) (*remdups-adj (sort ps)*))

lemma *set-prime-factorization-nat'*:

set (prime-factorization-nat' n) = ($\lambda p. (p, \text{multiplicity } p \ n - 1)$) ‘*prime-factors n*
<proof>

lemma *distinct-prime-factorization-nat' [simp]*: *distinct (prime-factorization-nat' n)*

<proof>

lemmas (**in** *multiplicative-function'*) *efficient-code'* =

efficient-code [of $\lambda-. \text{prime-factorization-nat' } n \ n$ **for** *n*,
 OF set-prime-factorization-nat' distinct-prime-factorization-nat']

14.1 Möbius μ function

definition *moebius-mu-aux* :: *nat* \Rightarrow (*unit* \Rightarrow *nat list*) \Rightarrow *int* **where**

moebius-mu-aux n ps =
 (*if* $n \neq 0 \wedge \neg 4 \text{ dvd } n \wedge \neg 9 \text{ dvd } n$ *then*
 (*let ps* = *ps* ()) *in if* *distinct ps* *then if* *even (length ps)* *then 1* *else -1* *else 0*) *else 0*)

lemma *moebius-mu-conv-moebius-mu-aux*:

fixes *qs* :: *unit* \Rightarrow *nat list*

defines *ps* \equiv *qs* ()

assumes *mset ps* = *prime-factorization n*

shows *moebius-mu n* = *of-int (moebius-mu-aux n qs)*

<proof>

lemma *moebius-mu-code* [code]:
 $moebius-mu\ n = of-int\ (moebius-mu-aux\ n\ (\lambda-. prime-factorization-nat\ n))$
 ⟨proof⟩

value *moebius-mu 12578972695257* :: int

14.2 Euler's ϕ function

primrec *totient-aux1* :: nat \Rightarrow nat list \Rightarrow nat **where**
 $totient-aux1\ n\ [] = n$
 $| totient-aux1\ n\ (p\ \#\ ps) = totient-aux1\ (n - n\ div\ p)\ ps$

lemma *of-nat-totient-aux1*:
assumes $\bigwedge p. p \in set\ ps \implies prime\ p \wedge p. p \in set\ ps \implies p\ dvd\ n\ distinct\ ps$
shows $real\ (totient-aux1\ n\ ps) = real\ n * (\prod_{p \in set\ ps} 1 - 1 / real\ p)$
 ⟨proof⟩

lemma *totient-conv-totient-aux1*:
assumes $set\ ps = prime-factors\ n\ distinct\ ps$
shows $totient\ n = totient-aux1\ n\ ps$
 ⟨proof⟩

definition *prime-factors-nat* :: nat \Rightarrow nat list **where**
 $prime-factors-nat\ n = remdups-adj\ (sort\ (prime-factorization-nat\ n))$

lemma *set-prime-factors-nat* [simp]: $set\ (prime-factors-nat\ n) = prime-factors\ n$
 ⟨proof⟩

lemma *distinct-prime-factors-nat* [simp]: $distinct\ (prime-factors-nat\ n)$
 ⟨proof⟩

definition *totient-aux2* :: (nat \times nat) list \Rightarrow nat **where**
 $totient-aux2\ xs = (\prod_{(p,k) \leftarrow xs} p^k * (p - 1))$

lemma *totient-conv-totient-aux2*:
assumes $n \neq 0$
assumes $set\ xs = (\lambda p. (p, multiplicity\ p\ n - 1))\ ` prime-factors\ n$
assumes $distinct\ xs$
shows $totient\ n = totient-aux2\ xs$
 ⟨proof⟩

lemma *totient-code1*: $totient\ n = totient-aux1\ n\ (prime-factors-nat\ n)$
 ⟨proof⟩

lemma *totient-code2*: $totient\ n = (if\ n = 0\ then\ 0\ else\ totient-aux2\ (prime-factorization-nat'\ n))$
 ⟨proof⟩

```

declare totient-code-naive [code del]

lemmas [code] = totient-code2

value totient 125789726827482323235784

```

14.3 Divisor Functions

```

lemmas [code del] = divisor-count-naive divisor-sum-naive
lemmas [code] = divisor-count.efficient-code' divisor-sum.efficient-code'

value int (divisor-count 378568418621)
value int (divisor-sum 378568418621)

```

14.4 Liouville's λ function

```

lemma [code]: liouville-lambda n =
  (if n = 0 then 0 else if even (length (prime-factorization-nat n)) then 1 else -1)
  ⟨proof⟩

value liouville-lambda 1264785343674 :: int

end

```

References

- [1] T. M. Apostol. *Introduction to Analytic Number Theory*. Undergraduate Texts in Mathematics. Springer-Verlag, 1976.