# Dedekind Sums

Manuel Eberl, Anthony Bordg, Lawrence C. Paulson, Wenda Li

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#### Abstract

For integers h, k, the Dedekind sum is defined as

$$s(h,k) = \sum_{r=1}^{k-1} \frac{r}{k} \left( \left\{ \frac{hr}{k} \right\} - \frac{1}{2} \right)$$

where  $\{x\} = x - \lfloor x \rfloor$  denotes the fractional part of x.

These sums occur in various contexts in analytic number theory, e.g. in the functional equation of the Dedekind  $\eta$  function or in the study of modular forms.

We give the definition of s(h,k) and prove its basic properties, including the reciprocity law

$$s(h,k) + s(k,h) = \frac{1}{12hk} + \frac{h}{12k} + \frac{k}{12h} - \frac{1}{4}$$

and various congruence results.

Our formalisation follows Chapter 3 of Apostol's  $Modular\ Functions$  and  $Dirichlet\ Series\ in\ Number\ Theory\ [1]$  and contains all facts related to Dedekind sums from it (without the exercises).

### Contents

1	Dedekind sums		
	1.1	Preliminaries	2
	1.2	Definition and basic properties	2
	1.3	The Reciprocity Law	9
	1.4	Congruence Properties	15

#### 1 Dedekind sums

notation dedekind\_frac (" $\langle \_ \rangle$ ")

show " $\langle x + 1 \rangle = \langle x \rangle$ " for x

proof

qed

```
theory Dedekind_Sums
imports
  Complex_Main
  "HOL-Library.Periodic_Fun"
  "HOL-Library.Real_Mod"
  "HOL-Number_Theory.Number_Theory"
  "Bernoulli_FPS"
begin
1.1
     Preliminaries
lemma rcong_of_intI: "[a = b] (mod m) \improx [of_int a = of_int b] (rmod
(of_int m))"
 by (metis cong_iff_lin mult.commute of_int_add of_int_mult rcong_altdef)
lemma rcong_of_int_iff: "[of_int a = of_int b] (rmod (of_int m)) ←→
[a = b] \pmod{m}"
proof
  assume "[of_int a = of_int b] (rmod (of_int m))"
  then obtain k where "real_of_int b = of_int a + of_int k * of_int m"
    unfolding rcong_altdef by blast
  also have "... = of_int (a + k * m)"
    by simp
  finally have "b = a + k * m"
    by linarith
  thus "[a = b] \pmod{m}"
    by (simp add: cong_iff_lin)
qed (intro rcong_of_intI)
      Definition and basic properties
definition dedekind_sum :: "int \Rightarrow int \Rightarrow real" where
  "dedekind_sum h k \equiv \sum r \in \{1... k\}. (of_int r / of_int k * (frac (of_int
(h*r) / of_int k) - 1/2))"
definition dedekind_frac :: "real ⇒ real" where
  "dedekind_frac x = (if x \in \mathbb{Z} then 0 else frac x - 1 / 2)"
lemma\ dedekind\_frac\_int\ [simp]:\ "x\ \in\ \mathbb{Z} \implies dedekind\_frac\ x\ =\ 0"
 by (auto simp: dedekind_frac_def)
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interpretation dedekind\_frac: periodic\_fun\_simple' dedekind\_frac

by (auto simp: dedekind\_frac\_def frac\_1\_eq)

```
lemma dedekind_frac_uminus [simp]: "\langle -x \rangle = -\langle x \rangle"
  by (auto simp add: dedekind_frac_def frac_neg)
lemma dedekind_frac_one_minus [simp]: "\langle 1 - x \rangle = -\langle x \rangle"
  by (metis dedekind_frac.minus_1 dedekind_frac_uminus minus_diff_eq)
lemma dedekind_frac_rcong:
  assumes "[x = x'] (rmod 1)"
           ''\langle x\rangle = \langle x'\rangle''
  \mathbf{shows}
proof -
  from assms obtain k where k: "x' = x + of_int k"
    by (auto simp: rcong_altdef)
  thus ?thesis
    by (auto simp: dedekind_frac.plus_of_int)
qed
lemma dedekind_frac_mod:
  "\langle of_{int} (a mod k) / of_{int} k \rangle = \langle of_{int} a / of_{int} k \rangle"
proof (cases "k = 0")
  case False
  have "[real_of_int (a mod k) = real_of_int a] (rmod real_of_int k)"
    by (intro rcong_of_intI cong_mod_leftI cong_refl)
  thus ?thesis using False
    by (intro dedekind_frac_rcong rcong_divide_modulus) auto
qed auto
lemma sum_dedekind_frac_eq_0:
  "(\sum r \in \{1... \langle k\}. \langle of_int \ r \ / \ of_int \ k \rangle) = 0"
proof -
  have "(\sum r \in \{1... < k\}. \langle of_int(r) / of_int(k)) =
          (\sum r \in \{1... < k\}. \langle of_int (k - r) / of_int k \rangle)"
    by (intro sum.reindex_bij_witness[of _ "\lambdar. k - r" "\lambdar. k - r"])
auto
  also have "(\sum r \in \{1... < k\}. \langle of_int(k-r) / of_int(k)) =
                (\sum r \in \{1... < k\}. \langle 1 + -of_int r / of_int k \rangle)"
    by (intro sum.cong refl arg_cong[of _ _ dedekind_frac]) (auto simp:
field_simps)
  also have "... = -(\sum r \in \{1... < k\}. \langle of_int r / of_int k \rangle)"
    by (simp add: sum_negf)
  finally show ?thesis
    by simp
qed
lemma sum_dedekind_aux:
  assumes "f (0::int) = 0"
  shows "(\sum r \in \{0... < k\}. f r) = (\sum r \in \{1... < k\}. f r)"
proof (rule sum.mono_neutral_right)
  have "\{0... < k\} - \{1... < k\} \subseteq \{0\}"
```

```
by auto
  thus "\forall i \in \{0... < k\} - \{1... < k\}. f i = 0"
     using assms by blast
qed auto
lemma sum_dedekind_frac_eq_0':
  "(\sum r \in \{0... < k\}. \langle of_int \ r \ / \ of_int \ k \rangle) = 0"
  have "(\sum r \in \{0... < k\}. \langle of_int r / of_int k \rangle) = (\sum r \in \{1... < k\}. \langle of_int k \rangle)
r / of_int k)"
     by (rule sum_dedekind_aux) auto
  with sum_dedekind_frac_eq_0[of k] show ?thesis
     by simp
qed
lemma sum dedekind frac mult eq 0:
  assumes "coprime h k"
            "(\sum r \in \{1..\langle k\}\}. \langle of_{int} (h * r) / of_{int} k \rangle) = 0"
  shows
  have "(\sum r \in \{1... < k\}. \langle of_int (h * r) / of_int k \rangle) =
          (\sum r \in \{1... < k\}. \langle of\_int ((h * r) mod k) / of\_int k \rangle)"
  proof (intro sum.cong)
     fix r assume r: "r \in \{1...< k\}"
     have k: "k > 0"
       using r by auto
     have "(h * r) \mod k = h * r - k * ((h * r) \operatorname{div} k)"
       by (metis minus_mult_div_eq_mod)
     also have "real_of_int ... / of_int k = of_int (h * r) / of_int k
- of_int ((h * r) div k)"
       using k by (auto simp: field_simps)
     also have "\langle ... \rangle = \langle of_{int} (h * r) / of_{int} k \rangle"
       by (rule dedekind_frac.minus_of_int)
     finally show "\langle of_{int} (h * r) / of_{int} k \rangle = \langle of_{int} ((h * r) mod k) \rangle
/ of_int k " ...
  qed auto
  also have "... = (\sum r \in \{1... < k\}. \langle of_int r / of_int k \rangle)"
     by (rule sum.reindex_bij_betw[OF bij_betw_int_remainders_mult[OF assms]])
  also have "... = 0"
     by (rule sum_dedekind_frac_eq_0)
  finally show ?thesis .
\mathbf{qed}
lemma sum_dedekind_frac_mult_eq_0':
  assumes "coprime h k"
            "(\sum r \in \{0... \langle k\}. \langle of_{int} (h * r) / of_{int} k \rangle) = 0"
  have "(\sum r \in \{0... < k\}. \langle of_{int} (h * r) / of_{int} k \rangle) = (\sum r \in \{1... < k\}. \langle of_{int} k \rangle)
(h * r) / of_int k)"
    by (intro sum_dedekind_aux) auto
```

```
also have "... = 0"
         by (intro sum_dedekind_frac_mult_eq_0 assms)
    finally show ?thesis .
lemma dedekind_sum_altdef:
    assumes "coprime h k"
                         "dedekind_sum h k = (\sum r \in \{1... < k\}). (of_int r / of_int k) * (of_int k)
 (h*r) / of_int k)"
proof -
    have "(\sum r \in \{1..\langle k\}\}. \langle of_{int} r / of_{int} k \rangle * \langle of_{int} (h*r) / of_{int}
                    (\sum r \in \{1..< k\}. (of_int r / of_int k - 1 / 2) * (of_int (h*r) / (h*
of_int k\rangle)"
    proof (intro sum.cong ballI)
         fix r assume r: "r \in \{1... < k\}"
         hence "1 \leq r" "r < k"
              by auto
         hence "\neg k dvd r"
              using zdvd_not_zless by auto
         hence "real_of_int r / real_of_int k \notin \mathbb{Z}"
              using r by (subst of_int_div_of_int_in_Ints_iff) auto
          moreover have "frac (real_of_int r / real_of_int k) = real_of_int
r / real_of_int k"
              using r by (subst frac_eq) auto
          ultimately show "\langle real\_of\_int r / real\_of\_int k \rangle * \langle real\_of\_int (h)
* r) / real_of_int k > =
                                                    (real_of_int r / real_of_int k - 1 / 2) * \( real_of_int \)
 (h * r) / real_of_int k"
              by (subst dedekind_frac_def) auto
    qed auto
    also have "... = (\sum r \in \{1... < k\}). of_int r / of_int k * (of_int (h*r) / (h*r))
of_int k\rangle) -
                                            1 / 2 * (\sum r \in \{1... < k\}. \langle of_int(h*r) / of_int(k))"
         unfolding sum_distrib_left sum_subtractf [symmetric] by (simp add:
ring distribs)
    also have "(\sum r \in \{1... k\}. \langle of_{int} (h*r) / of_{int} k \rangle) = 0"
          by (intro sum_dedekind_frac_mult_eq_0 assms)
    also have "(\sum r \in \{1... k\}. \text{ of_int } r \text{ / of_int } k * \text{ of_int } (h*r) \text{ / of_int }
k\rangle) = dedekind_sum h k"
         unfolding dedekind_sum_def
    proof (intro sum.cong)
         fix r assume r: "r \in \{1..\langle k\}"
         have "\neg k dvd h * r"
         proof
              assume "k \ dvd \ h * r"
              with assms have "k dvd r"
                   using coprime_commute coprime_dvd_mult_right_iff by blast
              hence "k \le r"
```

```
using r by (intro zdvd_imp_le) auto
                thus False
                     using r by simp
          qed
          hence "real_of_int (h * r) / real_of_int k ∉ Z"
                using r by (subst of_int_div_of_int_in_Ints_iff) auto
          thus "real_of_int r / real_of_int k * \( \text{real_of_int (h * r) / real_of_int } \)
k\rangle =
                           real\_of\_int \ r \ / \ real\_of\_int \ k \ * \ (frac \ (real\_of\_int \ (h \ * \ r) \ / \ )
real_of_int k) - 1 / 2)"
                by (auto simp: dedekind_frac_def)
     qed auto
     finally show ?thesis
          by simp
qed
theorem dedekind_sum_cong:
     assumes "[h' = h] \pmod{k}"
     assumes "coprime h' k \lor coprime h k"
     shows "dedekind_sum h' k = dedekind_sum h k"
proof -
     have coprime: "coprime h' k" "coprime h k"
           using coprime_cong_cong_left[OF assms(1)] assms(2) by auto
     show ?thesis
          unfolding coprime[THEN dedekind_sum_altdef]
     proof (intro sum.cong)
          fix r assume r: "r \in \{1..\langle k\}"
          have "[real_of_int (h' * r) = real_of_int (h * r)] (rmod real_of_int
k)"
                by (intro rcong_of_intI cong_mult cong_refl assms)
          hence "\langle real\_of\_int (h' * r) / real\_of\_int k \rangle = \langle real\_of\_int (h * rea
r) / real_of_int k "
                using r by (intro dedekind_frac_rcong rcong_divide_modulus) auto
          thus "(real_of_int r / real_of_int k) * (real_of_int (h' * r) / real_of_int
k\rangle =
                           ⟨real_of_int r / real_of_int k⟩ * ⟨real_of_int (h * r) / real_of_int
k\rangle"
                by simp
     qed auto
qed
theorem dedekind_sum_negate:
     assumes "coprime h k"
                            "dedekind_sum (-h) k = -dedekind_sum h k"
     \mathbf{shows}
proof -
     have *: "coprime (-h) k"
          using assms by auto
     show ?thesis
```

```
unfolding dedekind_sum_altdef[OF assms] dedekind_sum_altdef[OF *]
    by (simp_all add: sum_negf)
qed
theorem dedekind_sum_negate_cong:
  assumes "[h' = -h] (mod k)" "coprime h' k \lor coprime h k"
 shows "dedekind_sum h' k = -dedekind_sum h k"
  have coprime: "coprime h' k" "coprime h k"
    using coprime_cong_cong_left[OF assms(1)] assms(2) by auto
  from coprime show ?thesis
    using assms(1) dedekind_sum_cong dedekind_sum_negate by metis
qed
theorem dedekind_sum_inverse:
 assumes "[h * h' = 1] \pmod{k}"
 shows
           "dedekind_sum h k = dedekind_sum h' k"
proof -
 have 1: "coprime h k"
    using assms coprime_iff_invertible_int by blast
 have 2: "coprime h' k"
    using assms coprime_iff_invertible_int by (subst (asm) mult.commute)
  have "dedekind_sum h' k =
          (\sum r = 1..\langle k. \langle real\_of\_int (1 * r) / real\_of\_int k \rangle * \langle real\_of\_int
(h' * r) / real_of_int k)"
    unfolding dedekind_sum_altdef[OF 2] by simp
  also have "... = (\sum r = 1.. < k. \ (real\_of\_int \ (h * ((h' * r) mod k))) /
real_of_int k > *
                                 \real_of_int ((h' * r) mod k) / real_of_int
k\rangle)"
 proof (rule sum.cong, goal_cases)
    case (2 r)
    have "[real_of_int (1 * r) = real_of_int (h * h' * r)] (rmod real_of_int
k)"
      using assms by (intro roong of intI cong mult cong refl assms) (auto
simp: cong_sym)
    also have "[real_of_int (h * h' * r) = real_of_int (h * ((h' * r)
mod k))] (rmod of_int k)"
      by (subst mult.assoc, intro rcong_of_intI cong_mult cong_refl cong_mod_rightI)
    finally have *: "(real_of_int (1 * r) / real_of_int k) =
                      \real_of_int (h * ((h' * r) mod k)) / real_of_int
k\rangle"
      using 2 by (intro ext dedekind_frac_rcong rcong_divide_modulus)
auto
    have "[real_of_int (h' * r) = real_of_int (h' * r mod k)] (rmod real_of_int
k)"
      by (intro rcong_of_intI cong_refl cong_mod_rightI)
    hence "\(\real_of_int \)(h' * r) / real_of_int \(k\rangle = \real_of_int \)((h')
```

```
* r) mod k) / real_of_int k\"
      using 2 by (intro dedekind_frac_rcong rcong_divide_modulus) auto
    thus ?case using *
      by (simp add: mult_ac)
 ged auto
 also have "... = (\sum r = 1... < k. \ (real_of_int \ (h * r) / real_of_int \ k)
*  *  real_of_int r / real_of_int k)"
    by (rule sum.reindex_bij_betw [OF bij_betw_int_remainders_mult]) fact
 also have "... = dedekind_sum h \ k"
    using 1 by (simp add: dedekind_sum_altdef mult_ac)
 finally show ?thesis ..
qed
theorem dedekind_sum_inverse':
 assumes "[h * h' = -1] \pmod{k}"
 shows
          "dedekind_sum h k = -dedekind_sum h' k"
proof -
 have "[-(h * h') = - (-1)] \pmod{k}"
    using assms unfolding cong_minus_minus_iff .
  hence 1: "[h * -h' = 1] \pmod{k}"
    by simp
 hence "[h' * (-h) = 1] \pmod{k}"
    by (simp add: mult_ac)
 hence 2: "coprime h' k"
    using assms coprime_iff_invertible_int by blast
 from 1 have "dedekind_sum h k = dedekind_sum (-h') k"
    by (intro dedekind_sum_inverse)
  thus ?thesis
    using 2 by (simp add: dedekind_sum_negate)
qed
theorem dedekind_sum_eq_zero:
 assumes "[h^2 + 1 = 0] (mod k)"
 shows
         "dedekind_sum \ h \ k = 0"
proof -
 have "[h * h = h ^2 + 1 - 1] \pmod{k}"
    by (simp add: algebra_simps power2_eq_square)
 also have "[h ^2 + 1 - 1 = 0 - 1] (mod k)"
    by (intro cong_diff assms cong_refl)
 finally have "[h * h = -1] \pmod{k}"
    by simp
 hence "dedekind_sum h k = -dedekind_sum h k"
    by (rule dedekind_sum_inverse')
  thus ?thesis
    by simp
qed
```

#### 1.3 The Reciprocity Law

```
theorem sum_of_powers':
      "(\sum k < n :: nat. (real k) ^ m) = (bernpoly (Suc m) n - bernpoly (Suc m)
0) / (m + 1)"
proof (cases "n = 0")
     case True
     thus ?thesis
          by auto
next
     case False
     hence \{...,n\} = \{...,n-1\}
          by auto
     thus ?thesis
           using sum_of_powers[of m "n - 1"] False by (simp add: of_nat_diff)
qed
theorem sum_of_powers'_int:
     assumes "n \ge 0"
                         "(\sum k=0... < n::int. real_of_int k ^ m) = (bernpoly (Suc m) n
- bernpoly (Suc m) 0) / (m + 1)"
proof -
     have "(\sum k=0... < n::int. real_of_int k ^ m) = (\sum k < nat n::nat. real k
 ^ m)"
          by (intro sum.reindex_bij_witness[of _ int nat]) auto
     also have "(\sum k \le nat \ n :: nat. \ real \ k \ m) = (bernpoly (Suc \ m) \ n - bernpoly)
 (Suc m) 0) / (m + 1)"
          using assms by (subst sum_of_powers') (auto)
     finally show ?thesis by simp
qed
theorem sum_of_powers'_int_from_1:
     assumes "n \ge 0" "m > 0"
     shows "(\sum k=1.. < n::int. real_of_int k ^ m) = (bernpoly (Suc m) n
- bernpoly (Suc m) 0) / (m + 1)"
proof -
     have "(\sum k=1... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int k ^ m) = (\sum k=0... < n::int. real\_of\_int. real\_o
k ^ m)"
          using assms by (intro sum.mono_neutral_left) auto
     also have "... = (bernpoly (Suc m) n - bernpoly (Suc m) 0) / (m + 1)"
          using assms by (subst sum_of_powers'_int) auto
     finally show ?thesis by simp
qed
theorem dedekind_sum_reciprocity:
     assumes "h > 0" and "k > 0" and "coprime h k"
     shows "12 * h * k * dedekind_sum h k + 12 * k * h * dedekind_sum k h
                       h^2 + k^2 - 3 * h * k + 1"
```

```
proof -
  have upto_3_{eq} [simp]: "{...(3::nat)} = {0, 1, 2, 3}"
    by auto
  have [simp]: "(3 choose 2) = 3"
    by (simp add: eval_nat_numeral)
  define S where "S = (\sum r \in \{1... < k\}. \langle of\_int (h*r) / of\_int k \rangle ^ 2)" define T where "T = (\sum r = 1... < k. \lfloor real\_of\_int (h*r) / k \rfloor * (\lfloor real\_of\_int (h*r) / k \rfloor )"
(h * r) / k | + 1))"
  have "S = (\sum r \in \{1... k\}. \langle of_int((h * r) mod k) / of_int k \rangle ^ 2)"
    unfolding S_def
    by (intro sum.cong arg_cong[of _ _ " \lambda x.\ x \ ^2 "] dedekind_frac_mod
[symmetric] refl)
  also have "... = (\sum r \in \{1... < k\}. \langle of_int r / of_int k \rangle ^ 2)"
    by (rule sum.reindex_bij_betw[OF bij_betw_int_remainders_mult]) fact
  also have "... = (\sum r \in \{1... < k\}). (of_int r / of_int k - 1 / 2) ^ 2)"
  proof (rule sum.cong, goal_cases)
    case (2 r)
    have "\neg k dvd r"
       using 2 by (auto dest!: zdvd_imp_le)
    hence "of_int r / real_of_int k \notin \mathbb{Z}"
       using 2 by (subst of_int_div_of_int_in_Ints_iff) auto
    moreover have "frac (of_int r / real_of_int k) = of_int r / of_int
k"
       using 2 by (subst frac_eq) auto
    ultimately show ?case
       by (simp add: dedekind_frac_def)
  qed auto
  also have "... = 1 / k^2 * (\sum r=1... < k. of_int r^2) -
                     1 / k * (\sum_{r=1... < k. of_{int}} r) + (\sum_{r=1... < k.} 1 / 4)"
    unfolding sum_distrib_left sum_subtractf [symmetric] sum.distrib [symmetric]
dedekind_sum_def
    by (intro sum.cong) (auto simp: field_simps power2_eq_square)
  finally have "S = \dots ".
  note this [symmetric]
  also have "S = (\sum r = 1... < k.) (frac (real_of_int (h * r) / real_of_int
k) - 1 / 2)^2)"
    unfolding S_def
  proof (rule sum.cong, goal_cases)
    case (2 r)
    have "\neg k dvd r"
       using 2 by (auto dest!: zdvd_imp_le)
    hence "\neg k dvd h * r"
       using assms(3) coprime_commute coprime_dvd_mult_right_iff by blast
    hence "of_int (h * r) / real_of_int k \notin \mathbb{Z}"
       using 2 by (subst of_int_div_of_int_in_Ints_iff) auto
```

```
thus ?case
             by (simp add: dedekind_frac_def)
    qed auto
    also have "... =
         2 * h * dedekind_sum h k +
         \left(\sum r=1... < k. real\_of\_int \left(\left\lfloor of\_int(h*r)/k \right\rfloor * \left(\left\lfloor h*r/k \right\rfloor + 1\right)\right)\right) - \left(\left\lfloor h*r/k \right\rfloor + 1\right)\right)
         h^{2}/k^{2} * (\sum r=1... < k. real_of_int r^{2}) +
         (\sum r=1... < k. 1/4)"
         unfolding sum_distrib_left sum_subtractf [symmetric] sum.distrib [symmetric]
dedekind_sum_def
         by (intro sum.cong) (auto simp: field_simps power2_eq_square frac_def)
    finally have "2 * h * dedekind_sum h k =
                                       -(\sum r=1..< k. real\_of\_int ([of\_int(h*r)/k] * ([h*r/k])
+ 1))) +
                                        (h^2 + 1)/k^2 * (\sum r=1... < k. real of int r^2) - 1/k *
(\sum r=1... < k. of_int r)"
         by (simp add: add_divide_distrib ring_distribs)
    hence "6 * k * (2 * h * dedekind_sum h k) = 6 * k * (
                       -(\sum r=1... < k. real\_of\_int (|of\_int(h*r)/k| * (|h*r/k| + 1)))
                         (h^2 + 1)/k^2 * (\sum r=1.. < k. real\_of\_int r ^ 2) - 1/k * (\sum r=1.. < k.
of_int r))"
         by (simp only: )
    also have "... = -6 * k * (\sum r=1... < k. real_of_int (\lfloor of_int(h*r)/k \rfloor *
(|h*r/k| + 1))) +
                                          6 * (h^2 + 1)/k * (\sum r=1... < k. real_of_int r^2) -
                                          6 * (\sum r=1... < k. real_of_int r)"
         using \langle k \rangle 0> by (simp add: field_simps power2_eq_square)
    also have "(\sum v=1..\langle k. real\_of\_int v) = k^2/2 - k/2"
         using sum_of_powers'_int_from_1[of k 1] assms
         by (simp add: bernpoly_def algebra_simps power2_eq_square)
    also have "6 * ... = 3 * k^2 - 3 * k"
         using assms by (simp add: field_simps)
    also have "(\sum v=1... < k. real\_of\_int v ^ 2) = k ^ 3 / 3 - k ^ 2 / 2 +
         using \ assms \ by \ (subst \ sum\_of\_powers'\_int\_from\_1) \ (auto \ simp: \ bernpoly\_def)
    also have "real_of_int (6 * (h^2 + 1)) / real_of_int k * ... =
                              2 * h^{2} * k^{2} + 2 * k^{2} - 3 * h^{2} * k - 3 * k + h^{2} + 1"
         using assms by (simp add: field_simps power3_eq_cube power2_eq_square)
    finally have sum_eq1: "12 * h * k * dedekind_sum h k =
         -6 * k * T + 2 * h^2 * k^2 + 2 * k^2 - 3 * h^2 * k - 3 * k + h^2 + 1 -
3 * k^2 + 3 * k''
         by (simp add: mult_ac T_def)
    define N where "N = (\lambda v. \text{ card } \{r \in \{1... < k\}. | \text{real\_of\_int } (h*r)/k | = v
    have N_{eq} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\} = x \cdot |\{r \in \{1... < k\}\}. | real_of_int (h*r)/k | = v - 1\}
                                          \{1... < k\} \cap \{|k * (v - 1) / h| < ... | k * v / h|\}" for v
    proof -
```

```
have "[real\_of\_int (h*r)/k] = v - 1 \longleftrightarrow r \in \{[k * (v - 1) / h] < ... [k]\}
* v / h|}"
      if r: "r \in \{1..\langle k\}" \text{ for } r
    proof -
       have neg: "real_of_int r \neq k * v / h" for v
       proof
         assume "real_of_int r = k * v / h"
         hence "real_of_int (r * h) = real_of_int (k * v)"
           using assms unfolding of_int_mult by (simp add: field_simps)
         hence *: "r * h = k * v"
           by linarith
         have "k \ dvd \ r * h"
           by (subst *) auto
         with <coprime h k> have "k dvd r"
           by (simp add: coprime_commute coprime_dvd_mult_left_iff)
         with r show False
           by (auto dest!: zdvd imp le)
       qed
       have "|real_of_int (h*r)/k| = v - 1 \longleftrightarrow h * r / k \in {v-1..<v}"
         unfolding atLeastLessThan_iff by linarith
       also have "... \longleftrightarrow r \in \{k * (v - 1) / h.. < k * v / h\}"
         using assms by (auto simp: field_simps)
       also have "... \longleftrightarrow r \in \{k * (v - 1) / h < ... k * v / h\}"
         using neq[of "v - 1"] neq[of v] by auto
       also have "... \longleftrightarrow r \in \{|k * (v - 1) / h| < ... |k * v / h|\}"
         unfolding greaterThanAtMost_iff by safe linarith+
       finally show "|real_of_int (h*r)/k| = v - 1 \longleftrightarrow r \in {|k * (v -
1) / h | < ... | k * v / h | \}".
    qed
    thus "\{r \in \{1... < k\}. | real_of_int (h*r)/k | = v - 1\} =
            \{1... < k\} \cap \{|k * (v - 1) / h| < ... | k * v / h|\}"
       by blast
  qed
  define N' where "N' = (\lambda v. \text{ if } v = h \text{ then } k - 1 \text{ else } |k * v / h|)"
  have N eq: "int (N v) = N' v - N' (v - 1)" if v: "v \in \{1..h\}" for v
  proof (cases "v = h")
    case False
    have le: "[real_of_int k * (real_of_int v - 1) / real_of_int h]

  | real_of_int k * real_of_int v / real_of_int h | "
       using assms by (intro floor_mono divide_right_mono mult_left_mono)
auto
    have "N v = card ({1..<k} \cap {\lfloor k * (v - 1) / h \rfloor < .. | k * v / h | })"
       unfolding N_def N_eq_aux ..
    also have "\{|k * (v - 1) / h| < ... | k * v / h|\} \subseteq \{1... < k\}"
    proof -
       have ||k| * (v - 1) / h| \ge 0"
```

```
using v < k > 0 > < h > 0 >  by auto
       moreover have "\lfloor k * v / h \rfloor < k"
         using v False <k > 0> <h > 0> by (subst floor_less_iff) (auto
simp: field_simps)
       ultimately have "\{|k * (v - 1) / h| < ... | k * v / h|\} \subseteq \{0 < ... k-1\}"
         unfolding Ioc_subset_iff by auto
       also have "\{0<...k-1\} = \{1...< k\}"
         by force
       finally show ?thesis .
    hence "\{1..\langle k\} \cap \{\lfloor k*(v-1)/h \rfloor \langle ... \lfloor k*v/h \rfloor \} = \{\lfloor k*(v-1)/h \rfloor \}
/ h | < ... | k * v / h | \}"
       by blast
    finally show ?thesis
       using le v False by (simp add: N'_def)
    case [simp]: True
    have le: "|real_of_int k * (real_of_int h - 1) / real_of_int h | \le |
       using assms by (subst floor_le_iff) (auto simp: field_simps)
    have "N h = card (\{1..< k\} \cap \{|k * (h - 1) / h| < .. | k * h / h|\})"
       unfolding N_def N_eq_aux ..
    also have |k * h / h| = k
       using assms by simp
    also have \{1...< k\} \cap \{|\text{real\_of\_int } (k * (h - 1)) / \text{real\_of\_int } h|<...k\}
                 {|real_of_int (k * (h - 1)) / real_of_int h|+1..<k}"
    proof -
       have nonneg: "|real_of_int k * (real_of_int h - 1) / real_of_int
h \rfloor \geq 0"
         unfolding of_int_mult using assms by auto
       have \{1..\langle k\} \cap \{[real\_of\_int (k * (h - 1)) / real\_of\_int h] + 1..\langle k+1\}\}
              {\max 1 (\lceil real\_of\_int \ k * (real\_of\_int \ h - 1) / real\_of\_int}
h |+1)..<k}"
         by simp
       also have "max 1 (|real_of_int k * (real_of_int h - 1) / real_of_int
h \mid +1) =
                    |real_of_int k * (real_of_int h - 1) / real_of_int h |
+ 1"
         using nonneg by auto
       also have \{|real\_of\_int (k * (h - 1)) / real\_of\_int h| +1.. < k+1\}
                   {| real\_of\_int (k * (h - 1)) / real\_of\_int h | < ...k}"}
         by force
       finally show ?thesis by simp
    ged
    finally show ?thesis
       using le by (simp add: N'_def)
```

```
qed
```

```
have "T = (\sum (v,r) \in (SIGMA \ v:\{1..h\}. \ \{r\in\{1..\langle k\}. \ | of_int \ (h*r)/k | \})
= v - 1). (v - 1) * v"
    unfolding T_def
  proof (intro sum.reindex_bij_witness[of _ snd "\lambdar. (|of_int (h*r)/k|
+ 1, r)"], goal_cases)
    case (2 r)
    have "[real_of_int h * real_of_int r / real_of_int k] < h"
      using 2 assms by (subst floor_less_iff) (auto simp: field_simps)
    thus ?case using 2 assms by auto
  qed (use assms in auto)
 also have "... = (\sum v=1..h. \sum r/r \in \{1... < k\} \land \lfloor real\_of\_int(h*r)/k \rfloor = \{1... < k\}
v - 1. (v - 1) * v)"
  proof (intro sum.Sigma [symmetric] ballI)
    show "finite \{r \in \{1..\langle k\}\}. | real_of_int (h * r) / real_of_int k |
= v - 1 for v
      by (rule finite_subset[of _ "{1..<k}"]) auto</pre>
  qed auto
  also have "... = (\sum v=1..h. (v-1) * v * int (N v))"
    by (simp add: N_def mult_ac)
  also have "... = (\sum v=1..h. (v-1) * v * (N' v - N' (v-1)))"
    by (intro sum.cong) (auto simp: N_eq)
  also have "... = (\sum v=1..h. (v-1) * v * N' v) - (\sum v=1..h. (v-1)
* v * N' (v - 1))"
    by (simp add: ring_distribs sum_subtractf)
  also have "(\sum v=1..h. (v-1) * v * N' (v-1)) = (\sum v=0..< h. (v+1))
1) * v * N' v)"
    by (rule sum.reindex_bij_witness[of _ "\lambdav. v+1" "\lambdav. v-1"]) auto
  also have "... = (\sum v=1... < h. (v + 1) * v * N' v)"
    by (intro sum.mono_neutral_right) auto
  also have "{1..h} = insert h {1..<h}"
    using assms by auto
  also have "(\sum v \in .... (v - 1) * v * N' v) =
                (\sum v = 1.. < h. (v - 1) * v * N' v) + (h - 1) * h * N' h"
    by (subst sum.insert) auto
  also have "(\sum v=1... h. (v-1) * v * N' v) + (h-1) * h * N' h - (\sum v=1... h. h. h)
(v+1) * v * N' v) =
              (h-1) * h * N' h - 2 * (\sum v=1...< h. v * N' v)"
    by (simp add: ring_distribs sum_subtractf sum.distrib sum_distrib_left
sum_negf mult_ac)
  also have "(\sum v = 1... h. v * N' v) = (\sum v = 1... h. v * [of_int (k * l)])
v)/h|)"
    by (intro sum.cong) (auto simp: N'_def)
  also have "N' h = k - 1"
    by (simp add: N'_def)
  finally have *: "- 2 * (\sum v = 1... \langle h. v * [of_int (k * v)/h]) =
                    -(h - 1) * h * (k - 1) + T''
    by linarith
```

```
have "12 * k * h * dedekind_sum k h =
           6 * k * (-2 * (\sum v=1..\langle h. v * | real_of_int (k * v)/h|)) +
            12*k^2/h * (\sum v=1..< h. real_of_int v ^ 2) -
            6 * k * (\sum v=1...< h. real_of_int v)"
    by (simp add: dedekind_sum_def algebra_simps power2_eq_square sum.distrib
sum_subtractf
                  sum_distrib_left sum_distrib_right sum_negf frac_def
sum_divide_distrib)
  also note *
  also have "real_of_int (6 * k * (-(h - 1) * h * (k - 1) + T)) =
             6 * k * T - 6 * h^2 * k^2 + 6 * h * k^2 + 6 * h^2 * k - 6 * h
* k"
    by (simp add: algebra_simps power2_eq_square)
  also have "(\sum v=1..\langle h. real\_of\_int v) = h^2 / 2 - h / 2"
    using sum_of_powers'_int_from_1[of h 1] assms
    by (simp add: bernpoly_def algebra_simps power2_eq_square)
  also have "real_of_int (6 * k) * ... = 3 * h^2 * k - 3 * h * k"
    using assms by (simp add: field_simps)
  also have "(\sum v=1... h. real_of_int v ^ 2) = h ^ 3 / 3 - h ^ 2 / 2 +
    using assms by (subst sum_of_powers'_int_from_1) (auto simp: bernpoly_def)
  also have "real_of_int (12 * k<sup>2</sup>) / real_of_int h * ... =
             4 * h^2 * k^2 - 6 * h * k^2 + 2 * k^2"
    using assms by (simp add: field_simps power3_eq_cube power2_eq_square)
  finally have sum_eq2: "12 * k * h * dedekind_sum k h =
                           6 * k * T - 2 * h^2 * k^2 + 3 * h^2 * k - 3 * h
* k + 2 * k^2"
    by simp
  have "real_of_int (12 * h * k) * dedekind_sum h k + real_of_int (12
* k * h) * dedekind_sum k h =
        real_of_int (h^2 - 3 * h * k + k^2 + 1)"
    unfolding sum_eq1 sum_eq2 by simp
  thus ?thesis
    by simp
qed
theorem dedekind_sum_reciprocity':
  assumes "h > 0" and "k > 0" and "coprime h k"
 shows "dedekind_sum h k = -dedekind_sum k h + h / k / 12 + k / h / 12
- 1 / 4 + 1 / (12 * h * k)"
  using dedekind_sum_reciprocity[OF assms] assms
  by (auto simp: field_simps power2_eq_square)
```

#### 1.4 Congruence Properties

definition dedekind\_sum' :: "int  $\Rightarrow$  int  $\Rightarrow$  int" where

```
"dedekind_sum' h k = |6 * real_of_int k * dedekind_sum h k | "
lemma dedekind_sum'_cong:
    "[h = h'] (mod k) \Longrightarrow coprime h k \lor coprime h' k \Longrightarrow dedekind_sum'
h k = dedekind_sum' h' k"
    unfolding dedekind_sum'_def by (subst dedekind_sum_cong[of h h' k])
auto
lemma
    assumes "k > 0"
    shows
                       of_int_dedekind_sum':
                           "real_of_int (dedekind_sum' h k) = 6 * real_of_int k * dedekind_sum
h k"
                      dedekind_sum'_altdef:
        and
                           "dedekind_sum' h k = h * (k - 1) * (2 * k - 1) -
                                 6 * (\sum r = 1.. < k. r * [of_int (h * r) / k]) - 3 * (k * r)
(k - 1) div 2)"
                      dedekind_sum'_cong_3: "[dedekind_sum' h k = h * (k - 1) * (2)
* k - 1)] (mod 3)"
proof -
    have [simp]: "{..3} = {0,1,2,3::nat}"
    have [simp]: "(3 choose 2) = 3"
        by (auto simp: eval_nat_numeral)
    have "6 * k * dedekind_sum h k = 6 * h / k * (\sum r=1... < k. of_int r ^
2) -
                      6 * (\sum r=1... < k. of_int r * [h*r/k]) - 3 * (\sum r=1... < k. of_int)
r ^ 1)"
        by (simp add: dedekind_sum_def frac_def algebra_simps sum_distrib_left
sum_distrib_right
                                       sum_subtractf sum.distrib sum_divide_distrib power2_eq_square)
    also have "6 * h / k * (\sum r=1.. < k. real\_of\_int r ^ 2) = 2 * h * k^2 + k^2 
h - 3 * h * k"
        using assms by (subst sum_of_powers'_int_from_1)
                                          (auto simp: bernpoly_def field_simps power2_eq_square
power3_eq_cube)
    also have "3 * (\sum r=1.. < k. real_of_int r^1) = 3 * (k * (k - 1) / 2)"
        using assms by (subst sum_of_powers'_int_from_1)
                                          (auto simp: bernpoly_def field_simps power2_eq_square)
    also have "even (k * (k - 1))"
        by auto
    hence "(k * (k - 1) / 2) = real_of_int ((k * (k - 1)) div 2)"
        by fastforce
    also have "2 * h * k^2 + h - 3 * h * k = h * (k - 1) * (2 * k - 1)"
        by (simp add: algebra_simps power2_eq_square)
    finally have eq: "6 * real_of_int k * dedekind_sum h k =
                                           real_of_int (h * (k - 1) * (2 * k - 1) -
                                                  6 * (\sum r = 1... \langle k. r * [of_int (h * r) / k]) - 3
```

```
* (k * (k - 1) div 2))"
    unfolding of_int_diff of_int_add by simp
 have "6 * real_of_int k * dedekind_sum h k \in \mathbb{Z}"
    unfolding eq by (rule Ints_of_int)
  thus "real_of_int (dedekind_sum' h k) = 6 * of_int k * dedekind_sum
h k"
    unfolding dedekind_sum'_def by (auto elim!: Ints_cases)
 show eq': "dedekind_sum' h k = h * (k - 1) * (2 * k - 1) -
               6 * (\sum r = 1.. < k. r * [of_int (h * r) / k]) - 3 * (k * r)
(k - 1) div 2)"
    unfolding dedekind_sum'_def eq by (simp only: floor_of_int)
  have "[dedekind_sum' h k = h * (k - 1) * (2 * k - 1) - 0 - 0] (mod 3)"
    unfolding eq' by (intro cong_diff cong_refl) (auto simp: Cong.cong_def)
  thus "[dedekind sum' h k = h * (k - 1) * (2 * k - 1)] (mod 3)"
    by simp
qed
lemma three_dvd_dedekind_sum'_iff_aux:
  fixes h k :: int
  defines "\vartheta \equiv \gcd 3 k"
 assumes "k > 0" "coprime h k"
          "3 dvd (2 * dedekind_sum' h k) \longleftrightarrow \neg3 dvd k"
 shows
proof -
 have "[2 * dedekind_sum' h k = 2 * (h * (k - 1) * (2 * k - 1))] (mod
3)"
    by (intro cong_mult dedekind_sum'_cong_3 cong_refl assms)
 also have "2 * (h * (k - 1) * (2 * k - 1)) = h * (k - 1) * (4 * k + 1)
(-2))"
    by (simp add: algebra_simps)
  also have "[... = h * (k - 1) * (1 * k + 1)] (mod 3)"
    by (intro cong_mult cong_refl cong_add cong_diff) (auto simp: Cong.cong_def)
  finally have cong: [2 * dedekind_sum' h k = h * (k - 1) * (k + 1)] (mod
3)"
    by simp
 show "3 dvd (2 * dedekind_sum' h k) \longleftrightarrow \neg 3 dvd k"
  proof (cases "3 dvd k")
    case True
    have "¬3 dvd h"
      using <coprime h k> True by fastforce
    have "[2 * dedekind_sum' h k = h * (k - 1) * (k + 1)] (mod 3)"
      by (fact cong)
    also have "[h * (k - 1) * (k + 1) = h * (0 - 1) * (0 + 1)] (mod 3)"
      using True by (intro cong_mult cong_diff cong_add cong_refl) (auto
simp: cong_0_iff)
    also have "h * (0 - 1) * (0 + 1) = -h"
      by simp
    finally have "3 dvd (2 * dedekind_sum' h k) \longleftrightarrow 3 dvd (-h)"
```

```
using \ \textit{cong\_dvd\_iff} \ by \ \textit{blast}
    with \langle \neg 3 \text{ dvd } h \rangle True show ?thesis by auto
  next
    case False
    hence "3 dvd (k + 1) \lor 3 dvd (k - 1)"
      by presburger
    hence "3 dvd (h * (k - 1) * (k + 1))"
      by force
    also have "?this \longleftrightarrow 3 dvd (2 * dedekind_sum' h k)"
      using cong_cong_dvd_iff by blast
    finally show ?thesis
      using False by auto
  qed
qed
lemma dedekind_sum'_reciprocity:
  fixes h k :: int
  assumes "h > 0" "k > 0" "coprime h k"
  shows "2 * h * dedekind_sum' h k = -2 * k * dedekind_sum' k h + h<sup>2</sup>
+ k^2 - 3 * h * k + 1"
proof -
  have "real_of_int (2 * h * dedekind_sum' h k) = 12 * h * k * dedekind_sum
h k"
    unfolding of_int_mult of_int_dedekind_sum', [OF assms(2)] by simp
  also have "... = real_of_int (-12 * k * h) * dedekind_sum k h + real_of_int
(h^2 + k^2 - 3 * h * k + 1)"
    using dedekind_sum_reciprocity[OF assms] by simp
  also have "... = real_of_int (-2 * k * dedekind_sum' k h + h^2 + k^2 -
3 * h * k + 1)"
    using of_int_dedekind_sum', [OF assms(1)] by simp
  finally show ?thesis by linarith
qed
lemma cong_dedekind_sum'_1:
  fixes h k :: int
  defines "\vartheta \equiv \gcd 3 k"
  assumes "h > 0" "coprime h k"
  shows "[2 * k * dedekind_sum' k h = 0] (mod \vartheta * k)"
proof -
  have "\vartheta = 1 \vee \vartheta = 3"
    using gcd_prime_int[of 3 k] unfolding \theta_def by auto
  thus ?thesis
  proof
    assume "\vartheta = 3"
    hence "3 dvd k"
      unfolding \vartheta_{-}def by (metis gcd_dvd2)
    with <coprime h k> have "¬3 dvd h"
      by force
```

```
have "3 dvd 2 * dedekind_sum' k h"
      using assms <-3 dvd h>
      by (subst three_dvd_dedekind_sum'_iff_aux) (auto simp: coprime_commute)
    hence "3 * k dvd 2 * dedekind_sum' k h * k"
      by auto
    thus ?thesis
      using \langle \vartheta = 3 \rangle by (simp add: cong_0_iff)
  qed (auto simp: cong_0_iff)
qed
lemma cong_dedekind_sum'_2_aux:
 fixes h k :: int
  defines "\vartheta \equiv \gcd 3 k"
 assumes "h > 0" "k > 0" "coprime h k"
 shows "[2 * h * dedekind_sum' h k = h^2 + 1] (mod \vartheta * k)"
 have "[-(2 * k * dedekind_sum' k h) + h^2 + k^2 - 3 * h * k + 1 = -0 +
h^2 + 0 - 0 + 1] (mod \vartheta * k)"
   unfolding \vartheta_{-}def
    by (intro cong_diff cong_add cong_mult cong_refl cong_uminus cong_dedekind_sum'_1
assms)
       (simp_all add: cong_0_iff power2_eq_square)
 also have "-(2 * k * dedekind_sum' k h) + h^2 + k^2 - 3 * h * k + 1 =
2 * h * dedekind_sum' h k"
    using dedekind_sum'_reciprocity[of k h] assms by (auto simp: coprime_commute)
  finally show ?thesis by simp
qed
lemma dedekind_sum'_negate:
 assumes "k > 0" "coprime h k"
           "dedekind_sum' (-h) k = -dedekind_sum' h k"
 shows
proof -
 have "real_of_int (dedekind_sum' (-h) k) = real_of_int (-dedekind_sum'
h k)"
    using assms unfolding of_int_minus
    by (subst (1 2) of_int_dedekind_sum') (auto simp: dedekind_sum_negate)
 thus ?thesis
    by linarith
qed
lemma cong_dedekind_sum'_2:
  fixes h k :: int
  defines "\vartheta \equiv \gcd 3 k"
 assumes "k > 0" "coprime h k"
 shows "[2 * h * dedekind_sum' h k = h^2 + 1] (mod \vartheta * k)"
proof (cases h "0 :: int" rule: linorder_cases)
  case greater
  thus ?thesis
    using cong_dedekind_sum'_2_aux[of h k] assms by auto
```

```
next
  case less
  thus ?thesis
    using cong_dedekind_sum'_2_aux[of "-h" k] assms by (auto simp: dedekind_sum'_negate)
next
  case equal
  thus ?thesis using assms by (auto simp: Cong.cong_def)
qed
theorem dedekind_sum'_cong_8:
  assumes "k > 0" "coprime h k"
  shows "[2 * dedekind_sum' h k =
           (k-1)*(k+2) - 4*h*(k-1) + 4*(\sum r \in \{1...<(k+1) \ div \ 2\}. \ \lfloor 2*h*r/k \, \rfloor)]
(mod 8)"
proof -
  define S1 where "S1 = (\sum r=1... < k. r * \lfloor of_int (h * r) / k \rfloor)"
  define S2 where "S2 = (\sum r \mid r \in \{1... < k\} \land odd r. \mid of_int (h * r) /
  define S3 where "S3 = (\sum r=1...< k. \lfloor of_{int} (h * r) / k \rfloor)" define S4 where "S4 = (\sum r=1...< (k+1) div 2. \lfloor of_{int} (2 * h * r) / k \rfloor)"
  have "4 * 2 dvd 4 * (k * (k - 1))"
    by (intro mult_dvd_mono) auto
  hence dvd: "8 dvd 4 * k * (k - 1)"
    by (simp add: mult_ac)
  have "[4 * S1 = 4 * (\sum r=1... < k. if even r then 0 else of_int (h *
r) / k|)] (mod 8)"
    unfolding S1_def sum_distrib_left
  proof (intro cong_sum, goal_cases)
    case (1 r)
    show ?case
    proof (cases "odd r")
       case True
       have *: "4 * r \mod 8 = 4"
         using <odd r> by presburger
       have "[4 * r * | real_of_int (h * r) / k | = 4 * | real_of_int (h * r) / k |
r) / k | ] (mod 8) "
         by (rule cong_mult[OF\_cong\_refl]) (use * in <auto simp: Cong.cong\_def>)
       thus ?thesis
         using True by (simp add: mult_ac)
    qed (auto simp: cong_0_iff)
  also have "(\sum r=1..\langle k. if even r then 0 else \lfloor of_{int} (h * r) / k \rfloor) =
    unfolding S2_def by (intro sum.mono_neutral_cong_right) auto
  finally have S12: "[4 * S1 = 4 * S2] \pmod{8}".
```

```
have "2 * dedekind_sum' h k =
           2 * (h * (k - 1) * (2 * k - 1)) - 12 * S1 - 6 * (k * (k - 1)
div 2)"
    using assms by (subst dedekind_sum'_altdef) (auto simp: S1_def)
  also have "6 * (k * (k - 1) div 2) = 3 * k * (k - 1)"
    by (subst div_mult_swap) auto
  also have "2 * (h * (k - 1) * (2 * k - 1)) - 12 * S1 - 3 * k * (k -
              -2 * h * (k - 1) + h * (4 * k * (k - 1)) - 12 * S1 + k *
(k-1)-4*k*(k-1)"
    by (simp add: algebra_simps)
  also have "[... = -2 * h * (k - 1) + h * 0 - 4 * S1 + k * (k - 1) - 0]
(mod 8)"
    using dvd by (intro cong_diff cong_add S12 cong_mult cong_refl)
                   (auto simp: Cong.cong_def mod_eq_0_iff_dvd)
  also have "-2 * h * (k - 1) + h * 0 - 4 * S1 + k * (k - 1) - 0 = (k
-1) * (k - 2 * h) - 4 * S1"
    by (simp add: algebra_simps)
  also have "[... = (k - 1) * (k - 2 * h) - 4 * S2] (mod 8)"
    by (intro cong_diff cong_refl S12)
  also have "S2 = S3 - S4"
  proof -
    have *: "\{r. r \in \{1... \le k\} \land odd r\} = \{1... \le k\} - \{r. r \in \{1... \le k\} \land k\} 
even r}"
      by auto
    have "S2 = S3 - (\sum r \mid r \in \{1..\langle k\} \land \text{ even } r. \mid \text{of\_int } (h * r) / k \mid)"
      unfolding S2_def S3_def * by (subst Groups_Big.sum_diff) auto
    also have "(\sum r \mid r \in \{1... k\} \land \text{ even } r. \mid \text{of\_int } (h * r) \mid k)" = S4"
      unfolding S4_def using <k > 0>
      by (intro sum.reindex_bij_witness[of _ "\lambdar. 2 * r" "\lambdar. r div 2"])
          (auto simp: real_of_int_div)
    finally show ?thesis .
  qed
  also have "4 * (S3 - S4) = 4 * S3 - 4 * S4"
    by (simp add: algebra_simps)
  also have "4 * S3 = 2 * (h - 1) * (k - 1)"
  proof -
    have "real_of_int S3 = (\sum r=1..\langle k. -\langle of_int (h * r) / k \rangle + of_int
(h * r) / k - 1 / 2)"
      unfolding S3_def of_int_sum
    proof (intro sum.cong)
      fix r assume r: "r \in \{1..\langle k\}"
      hence "\neg k dvd r"
        by (auto dest!: zdvd_imp_le)
      hence "\neg k dvd (h * r)"
        using assms coprime_commute coprime_dvd_mult_right_iff by blast
      hence "real_of_int (h * r) / real_of_int k \notin \mathbb{Z}"
        using assms by (subst of_int_div_of_int_in_Ints_iff) auto
      thus "real_of_int |real_of_int (h * r) / real_of_int k =
```

```
-\langle real\_of\_int (h * r) / real\_of\_int k \rangle + real\_of\_int (h
* r) / real_of_int k - 1 / 2"
                  by (simp add: dedekind_frac_def frac_def)
         also have "4 * ... = -4 * (\sum r=1... \langle k... \langle of\_int (h * r) / k \rangle) +
                                                            4 * h / k * (\sum r=1.. < k. of_int r^1) - ((real_of_int r^1) - ((real_of_int r^2)) - ((r
k * 4 - 4) / 2)"
             using assms
             by (simp add: sum.distrib sum_subtractf sum_negf of_nat_diff mult_ac
                                              sum_distrib_left sum_distrib_right sum_divide_distrib)
         also have "(\sum r=1... \langle k. \langle of_int(h*r)/k \rangle) = 0"
             using assms by (intro sum_dedekind_frac_mult_eq_0)
         also have "4 * h / k * (\sum r=1... < k. real_of_int r ^ 1) = 2 * h * (k...)
- 1)"
             using assms by (subst sum_of_powers'_int_from_1) (auto simp: field_simps
bernpoly def)
        also have "-4 * 0 + real_of_int (2 * h * (k - 1)) - (real_of_int k
* 4 - 4) / 2 =
                                  real_of_int (2 * (h - 1) * (k - 1))"
             by (simp add: field_simps)
         also have "4 * real_of_int S3 = real_of_int (4 * S3)"
             by simp
         finally show ?thesis by linarith
    qed
    also have "(k - 1) * (k - 2 * h) - (2 * (h - 1) * (k - 1) - 4 * S4)
                              (k-1)*(k+2)-4*h*(k-1)+4*S4"
         by (simp add: algebra_simps)
    finally show "[2 * dedekind_sum' h k = (k-1)*(k+2) - 4*h*(k-1) + 4*S4]
(mod 8)"
         unfolding S4_def .
qed
theorem dedekind_sum'_cong_8_odd:
    assumes "k > 0" "coprime h k" "odd k"
    shows "[2 * dedekind_sum' h k =
                      k - 1 + 4*(\sum r \in \{1... < (k+1) \text{ div } 2\}. |2*h*r/k|)] \pmod{8}"
proof -
    define S where "S = (\sum r=1...<(k+1) \text{ div } 2. | \text{of_int } (2 * h * r) / k|)"
    have 1: [(k-1)*(k+2) = k - 1] \pmod{8}
    proof -
         from assms obtain k' where k': "k = 2 * k' + 1"
             by (elim oddE)
         define k'', where "k'' = k' \mod 4"
         have "k'', \in \{0...<4\}"
             unfolding k''_def by simp
         also have "{0..<4} = {0, 1, 2, 3::int}"
```

```
by auto
    finally have "(2 * k'' + 1) ^ 2 \mod 8 = 1"
      by auto
    have "[k ^2 = ((2 * k') \mod (2 * 4) + 1) ^2 \mod 8] \pmod 8"
      unfolding k' by (intro cong_add cong_refl cong_pow cong_mod_rightI)
(auto simp: Cong.cong_def)
    also have "(2 * k) mod (2 * 4) = 2 * k'"
      by (subst mod_mult_mult1) (auto simp: k''_def)
    also have "(2 * k'' + 1) ^ 2 \mod 8 = 1"
      by fact
    finally have *: "[k \hat{2} = 1] (mod 8)".
   have (k-1)*(k+2) = k^2 + k - 2
      by (simp add: algebra_simps power2_eq_square)
    also have "[... = 1 + k - 2] (mod 8)"
      by (intro cong_add cong_refl cong_diff *)
   finally show ?thesis
      by simp
  qed
 have 2: [4 * h * (k - 1) = 0] \pmod{8}
 proof -
   have "4 * 1 * 2 dvd 4 * h * (k - 1)"
      using assms by (intro mult_dvd_mono) auto
    thus ?thesis by (simp add: cong_0_iff)
  qed
 have "[2 * dedekind_sum' h k = (k-1)*(k+2) - 4*h*(k-1) + 4*S] (mod 8)"
    unfolding S_def using assms by (intro dedekind_sum'_cong_8)
  also have "[(k-1)*(k+2) - 4*h*(k-1) + 4*S = (k-1) - 0 + 4 * S] (mod
8)"
    by (intro cong_add cong_diff cong_refl 1 2)
 finally show ?thesis
    by (simp add: S_def)
qed
lemma dedekind_sum'_cong_power_of_two:
  fixes h \ k \ k1 :: int and n :: nat
 assumes "h > 0" "k1 > 0" "odd k1" "n > 0" "k = 2 \hat{n} * k1" "coprime
h k"
 shows
          "[2 * h * dedekind_sum' h k =
              h^2 + k^2 + 1 + 5 * k - 4 * k * (\sum v=1...<(h+1) div 2. [of_int]
(2 * k * v) / h|)]
           (mod 2 ^ (n + 3))"
proof -
 from assms have "even k"
```

```
by auto
    with <coprime h k> have "odd h"
        using coprime_common_divisor odd_one by blast
    from assms have "k > 0"
        by auto
    define S where "S = (\sum v=1...<(h+1) \ div \ 2. \ | of_int \ (2 * k * v) / h |)"
    have "[2 * dedekind_sum' k h * k = (h - 1 + 4 * S) * k] (mod (8 * k))"
        unfolding S_{def} using \langle h > 0 \rangle \langle k > 0 \rangle \langle coprime | h | k \rangle \langle odd | h \rangle
        by (intro dedekind_sum'_cong_8_odd cong_cmult_rightI) (auto simp:
coprime_commute)
    also have "8 * k = 2 ^ (n + 3) * k1"
        using assms by (simp add: power_add)
    finally have *: "[2 * dedekind_sum' k h * k = (h - 1 + 4 * S) * k] (mod
2 ^ (n + 3))"
        using cong_modulus_mult by blast
    have **: "[k - 4 * h * k = 5 * k] \pmod{2^{(n+3)}}"
    proof -
        have "2 ^ 2 * 2 ^ n * 2 dvd 4 * k * (h + 1)"
             using assms by (intro mult_dvd_mono) auto
        hence "2 \hat{} (n + 3) dvd (5 * k - (k - 4 * h * k))"
             by (simp add: algebra_simps power_add)
        thus ?thesis
             by (subst cong_sym) (auto simp: cong_iff_dvd_diff)
    have "2 * h * dedekind_sum' h k = h^2 + k^2 - 3 * h * k + 1 - 2 * dedekind_sum'
k h * k"
        using dedekind_sum'_reciprocity[of h k] <h > 0> <k > 0> <coprime h
k> by simp
    also have "[... = h^2 + k^2 - 3 * h * k + 1 - (h - 1 + 4 * S) * k] (mod
2 ^ (n + 3))"
        using * by (intro cong_add cong_diff cong_refl)
    also have h^2 + h^2 - 3hk + 1 - (h - 1 + 4k) + h = h^2 + h^2 + 1 + (h - 1) + h = h^2 + h^2 + h + h^2 + h = h^2 + h^2 +
- 4*h*k) - 4*k*S"
        by (simp add: algebra simps)
    also have "[... = h^2 + k^2 + 1 + 5 * k - 4*k*S] (mod 2 ^ (n + 3))"
        by (intro cong_add cong_diff cong_refl **)
    finally show ?thesis
        by (simp add: S_def)
qed
lemma dedekind_sum'_cong_power_of_two':
    fixes h k k1 :: int
    assumes "h > 0" "k > 0" "even k" "coprime h k"
                    "[2 * h * dedekind_sum' h k =
                              h^2 + k^2 + 1 + 5 * k - 4 * k * (\sum v=1..<(h+1) div 2. [of_int]
(2 * k * v) / h|)]
                         (mod \ 2 \ \hat{} \ (multiplicity \ 2 \ k \ + \ 3))"
```

```
{\bf proof} \ ({\it rule dedekind\_sum'\_cong\_power\_of\_two})
  define k1 where "k1 = k div 2 ^ multiplicity 2 k"
  have "2 ^ multiplicity 2 k dvd k"
    using multiplicity_dvd by blast
  thus k_{eq}: "k = 2 ^ multiplicity 2 k * k1"
    by (auto simp: k1_def)
  show "odd k1"
    using \langle k \rangle 0 \rangle multiplicity_decompose[of k 2] by (auto simp: k1_def)
  show "multiplicity 2 k > 0"
    using <even k> <k > 0> by (simp add: multiplicity_gt_zero_iff)
  show "k1 > 0"
    using \langle k \rangle 0 \rangle by (subst (asm) k_eq) (auto simp: zero_less_mult_iff)
qed (use assms in auto)
lemma dedekind_sum_diff_even_int_aux:
  fixes a b c d :: int assumes det: "a * d - b * c = 1"
  fixes q c1 r \delta' :: int and \delta :: real
  assumes a: "a > 0"
  assumes q: "q \in \{3, 5, 7, 13\}" and "c1 > 0"
  assumes c: "c = q * c1"
  defines "r \equiv 24 \text{ div } (q - 1)"
  defines "\delta' \equiv (2 * dedekind_sum' a c - (a + d)) - (2 * q * dedekind_sum'
a c1 - (a + d) * q)"
  defines "\delta \equiv (\text{dedekind\_sum a } c - (a+d)/(12*c)) - (\text{dedekind\_sum a } c1
- (a+d)/(12*c1))"
            "of_int \delta' = 12 * c * \delta" and "24 * c dvd r * \delta'"
  \mathbf{shows}
proof -
  define \vartheta where "\vartheta = gcd 3 c"
  define \vartheta 1 where "\vartheta 1 = gcd 3 c1"
  have "even r" "odd q"
    using q by (auto simp: r_def)
  have "q > 0"
    using q by auto
  have "c > 0"
    using q and \langle c1 \rangle 0> and c by auto
  have "[a * d - b * 0 = a * d - b * c] \pmod{c}"
    by (intro cong_diff cong_mult cong_refl) (auto simp: Cong.cong_def)
  also have "a * d - b * c = 1"
    by fact
  finally have "[a * d = 1] \pmod{c}"
    by simp
  hence "coprime a c"
    using coprime_iff_invertible_int by blast
  have "coprime a c1"
    using <coprime a c> c by auto
```

```
show of_int_\delta': "of_int \delta' = 12 * c * \delta"
    using \langle c \rangle 0> \langle c1 \rangle 0> \langle q \rangle 0> unfolding \delta'_def \delta_def of_int_mult
of_int_diff
    by (subst (1 2) of_int_dedekind_sum') (auto simp: field_simps c)
  have "[2 * a * dedekind_sum' a c = a^2 + 1] (mod \vartheta * c)"
    using <coprime a c> <c > 0> unfolding \vartheta_{def} by (intro cong_dedekind_sum'_2)
  hence "[2 * a * dedekind_sum' a c - a * (a + d) = (a^2 + 1) - a * (a + d)
+ d)] (mod \vartheta * c)"
    by (intro cong_diff cong_refl)
  also have "(a^2 + 1) - a * (a + d) = -b * c"
    using det unfolding power2_eq_square by (simp add: algebra_simps)
  finally have "[2 * a * dedekind_sum' a c - a * (a + d) = -b * c] (mod
\vartheta * c)".
  hence 1: "[2 * a * dedekind_sum' a c - a * (a + d) = -b * c] (mod \vartheta1
* c)"
    by (rule cong_dvd_mono_modulus) (auto simp: \vartheta_def \vartheta1_def c)
  have "[2 * a * dedekind_sum' a c1 * q = (a<sup>2</sup> + 1) * q] (mod \vartheta1 * c1
    using \langle coprime \ a \ c1 \rangle \langle c1 \rangle 0 \rangle unfolding \vartheta 1_{def}
    by (intro cong_cmult_rightI cong_dedekind_sum'_2) auto
  hence "[2 * a * dedekind_sum' a c1 * q - a * (a + d) * q =
           (a^2 + 1) * q - a * (a + d) * q] \pmod{\vartheta 1 * c}"
    by (intro cong_diff cong_refl) (auto simp: c mult_ac)
  also have "(a^2 + 1) * q - a * (a + d) * q = -q * b * c"
    using det unfolding power2_eq_square by (simp add: algebra_simps)
  finally have 2: "[2 * a * q * dedekind_sum' a c1 - a * (a + d) * q =
-q * b * c] (mod \vartheta1 * c)"
    by (simp add: mult_ac)
  have r_\delta: "r * \delta' = r * (2 * dedekind_sum' a c - (a + d) - c
                        (2 * q * dedekind_sum' a c1 - (a + d) * q))"
    by (simp add: \delta'_def algebra_simps)
  have r a \delta': "r * a * \delta' = r * (2 * a * dedekind sum' a c - a * (a +
d) -
                        (2 * a * q * dedekind_sum' a c1 - a * (a + d) * q))"
    by (simp add: \delta '_def algebra_simps)
  also have "[... = r * (-b * c - (-q * b * c))] (mod \vartheta 1 * c)"
    by (intro cong_mult cong_diff 1 2 cong_refl)
  also have "r * (-b * c - (-q * b * c)) = r * (q - 1) * b * c"
    by (simp add: algebra_simps)
  also have "r * (q - 1) = 24"
    using q by (auto simp: r_def)
  also have "\vartheta1 * 1 * c dvd 24 * b * c"
    by (intro mult_dvd_mono) (auto simp: \vartheta1_def gcd_dvdI1)
  hence "[24 * b * c = 0] (mod \vartheta1 * c)"
```

```
by (simp add: cong_0_iff)
  finally have "\vartheta1 * c dvd r * \delta' * a"
    by (simp add: cong_0_iff mult_ac)
  moreover have "coprime a (gcd 3 c1)"
    using <coprime a c1> coprime_imp_coprime dvd_trans by blast
  ultimately have "\vartheta1 * c dvd r * \delta'"
    using <coprime a c>
    by (subst (asm) coprime_dvd_mult_left_iff) (auto simp: \vartheta1_def coprime_commute)
  have "3 * c dvd r * \delta'"
  proof (cases "q = 3")
    case False
    hence "\neg 3 dvd q"
       using q by auto
    hence "coprime 3 q"
       by (intro prime_imp_coprime) auto
    hence [simp]: "\vartheta1 = \vartheta"
       by (auto simp: \theta_{def} \theta_{1_{def}} c \ gcd_{mult_right_left_cancel})
    show ?thesis
    proof (cases "\vartheta = 3")
       case True
       with \langle \vartheta 1 * c \text{ dvd } r * \delta' \rangle show ?thesis by auto
    \mathbf{next}
       case False
       hence [simp]: "\vartheta = 1"
         unfolding \vartheta_def by (subst gcd_prime_int) auto
       hence "\neg 3 dvd c"
         unfolding \theta_{\text{def}} by (subst (asm) gcd_prime_int) auto
       hence "¬3 dvd c1"
         unfolding c using <coprime 3 q> coprime_dvd_mult_right_iff by
blast
       have "[2 * dedekind_sum' a c = 0] (mod 3)"
         using three_dvd_dedekind_sum'_iff_aux[of c a] <c > 0> <coprime</pre>
a c > \langle \neg 3 \text{ dvd } c \rangle
         by (auto simp: cong_0_iff)
       moreover have "[2 * dedekind_sum' a c1 * q = 0] \pmod{3}"
         using three_dvd_dedekind_sum'_iff_aux[of c1 a] <c1 > 0> <coprime
a c1 \langle \neg 3 \ dvd \ c1 \rangle
         by (auto simp: cong_0_iff)
       ultimately have "[r * \delta]" =
                           r * ((0 - (a + d)) - (0 - (a + d) * q))] (mod
3)"
         unfolding r_{\delta} by (intro cong_mult cong_diff cong_refl) (simp_all
add: mult_ac)
       also have "r * ((0 - (a + d)) - (0 - (a + d) * q)) = r * (q - 1)
* (a + d)"
         by (simp add: algebra_simps)
       also have "r * (q - 1) = 24"
         using q by (auto simp: r_def)
```

```
also have "[24 * (a + d) = 0] \pmod{3}"
        by (simp add: cong_0_iff)
      finally have "3 dvd r * \delta,"
        by (simp add: cong_0_iff)
      moreover from \langle \vartheta 1 * c \ dvd \ r * \delta' \rangle have "c dvd r * \delta'"
        by (simp add: mult_ac)
      moreover have "coprime 3 c"
        using <-3 dvd c> by (intro prime_imp_coprime) auto
      ultimately show "3 * c dvd r * \delta'"
        using divides_mult by blast
    qed
  next
    case [simp]: True
    have [simp]: "r = 12"
      by (simp add: r def)
    have [simp]: "\vartheta = 3"
      by (auto simp: \theta_{def} c)
    show ?thesis
    proof (cases "\vartheta1 = 3")
      case True
      with \langle \vartheta 1 * c \text{ dvd } r * \delta' \rangle show ?thesis by simp
    next
      case False
      hence [simp]: "\vartheta1 = 1"
        unfolding \vartheta1_def by (subst gcd_prime_int) auto
      hence "¬3 dvd c1"
        unfolding \vartheta1_def by (subst (asm) gcd_prime_int) auto
      have "[2 * dedekind_sum' a c1 * a * q = 0 * q] (mod 3 * q)"
        using three_dvd_dedekind_sum'_iff_aux[of c1 a] <c1 > 0> <coprime
a c1 \langle \neg 3 \ dvd \ c1 \rangle
        by (intro cong_cmult_rightI) (auto simp: cong_0_iff)
      hence "[r * a * \delta' = r * (2 * a * dedekind_sum' a c - a * (a + d)]
-(0 - a * (a + d) * q))] (mod 9)"
        unfolding r_a_\delta by (intro cong_mult cong_diff cong_refl) (auto
simp: mult_ac)
      also have "r * (2 * a * dedekind sum' a c - a * (a + d) - (0 - a)
*(a + d) * q)) =
                  r * (2 * a * dedekind_sum' a c) + 2 * r * (a^2 + a * d)"
        by (simp add: algebra_simps power2_eq_square)
      also have "[2 * a * dedekind_sum' a c = a^2 + 1] (mod \vartheta * c)"
        using cong_dedekind_sum'_2[of c a] <c > 0> <coprime a c> unfold-
ing \theta_{-}def by auto
      hence "[2 * a * dedekind_sum' a c = a^2 + 1] (mod 9)"
        by (rule cong_dvd_mono_modulus) (auto simp: c)
      hence "[r * (2 * a * dedekind_sum' a c) + 2 * r * (a^2 + a * d) =
               r * (a^2 + 1) + 2 * r * (a^2 + a * d)] \pmod{9}"
        by (intro cong_add cong_mult cong_refl)
      also have "r * (a^2 + 1) + 2 * r * (a^2 + a * d) = 3 * r * a^2 + r
+ 2 * r * (a * d)"
```

```
by (simp add: algebra_simps)
      also have "a * d = b * c + 1"
         using det by (simp add: algebra_simps)
      also have "3 * r * a^2 + r + 2 * r * (b * c + 1) = 9 * (4 + 8 * b)
* c1 + a^2 * 4)"
         by (simp add: algebra_simps c)
      also have "[... = 0] (mod 9)"
         by (simp add: cong_0_iff)
      finally have "9 dvd r * a * \delta'"
         by (simp add: cong_0_iff)
      moreover have "coprime a (3 ^ 2)"
         using <coprime a c > c by (subst coprime_power_right_iff) auto
      hence "coprime a 9"
         by (simp del: coprime_power_right_iff)
      ultimately have "9 dvd r * \delta'"
         by (metis coprime_commute coprime_dvd_mult_right_iff mult.assoc
mult.commute)
      have "3 * c1 dvd 4 * \delta'"
      proof (rule divides_mult)
         show "coprime 3 c1"
           using \langle \vartheta 1 = 1 \rangle unfolding \vartheta 1_{def} by auto
         have "3 * c1 dvd 3 * (4 * \delta)"
           using \langle \vartheta 1 * c \ dvd \ r * \delta' \rangle by (simp add: c mult_ac)
         thus "c1 dvd 4 * \delta'"
           by (subst (asm) dvd_mult_cancel_left) auto
      next
         have "3 * 3 dvd 3 * (4 * \delta)"
           using <9 dvd r * \delta'> by simp
         thus "3 dvd 4 * \delta'"
           by (subst (asm) dvd_mult_cancel_left) auto
      hence "3 * c dvd 3 * (4 * \delta)"
         by (intro mult_dvd_mono dvd_refl) (auto simp: c)
      thus ?thesis
         by (simp add: c mult_ac)
    qed
  qed
  show "24 * c dvd r * \delta'"
  \operatorname{proof} (cases "even c")
    assume "odd c"
    hence "odd c1"
      by (auto simp: c)
    define T :: "int \Rightarrow int" where
       "T = (\lambda c. (\sum r = 1.. < (c + 1) \text{ div } 2. | \text{real\_of\_int } (2 * a * r) / \text{real\_of\_int}]
c|))"
```

```
have "[2 * dedekind_sum' a c = c - 1 + 4 * T c] (mod 8)"
      unfolding T_def by (rule dedekind_sum'_cong_8_odd)
                           (use \langle coprime \ a \ c \rangle \langle odd \ c \rangle \langle c \rangle \langle o \rangle  in auto)
    moreover have "[2 * dedekind_sum' a c1 * q = (c1 - 1 + 4 * T c1)
* a] (mod 8)"
      unfolding \ \textit{T\_def} \ by \ (\textit{intro cong\_mult cong\_refl dedekind\_sum'\_cong\_8\_odd})
                            (use \langle coprime \ a \ c1 \rangle \langle odd \ c1 \rangle \langle c1 \rangle \langle obs \ in \ auto)
    ultimately have "[r * \delta' = r * ((c - 1 + 4 * T c - (a + d)) -
                                      ((c1 - 1 + 4 * T c1) * q - (a + d)
* q))] (mod 8)"
      unfolding \delta'_def by (intro cong_diff cong_mult[of r] cong_refl)
(auto simp: mult_ac)
    also have "r * ((c - 1 + 4*T c - (a+d)) - ((c1 - 1 + 4*T c1)*q - (a+d)*q))
                  r * (q - 1) * (a + d + 1) + (r * 4) * (T c - q * T c1)"
      by (simp add: c algebra_simps)
    also have "r * (q - 1) = 24"
      using q by (auto simp: r_def)
    also have "[24 * (a + d + 1) + r * 4 * (T c - q * T c1) =
                 0 * (a + d + 1) + 0 * (T c - q * T c1)] (mod 8)"
      using <even r> by (intro cong_add cong_mult cong_refl) (auto simp:
cong_0_iff)
    finally have "8 dvd r * \delta,"
      by (simp add: cong_0_iff)
    have "8 * (3 * c) dvd r * \delta'"
    proof (rule divides_mult)
      have "coprime (2 \hat{3}) (3 * c)"
        using <odd c> unfolding coprime_power_left_iff by auto
      thus "coprime 8 (3 * c)"
        by (simp del: coprime_power_left_iff)
    qed fact+
    thus ?thesis
      by simp
 next
    assume "even c"
    with <coprime a c> have "odd a"
      using coprime_common_divisor odd_one by blast
    from <even c> and <odd q> have "even c1"
      by (auto simp: c)
    define n where "n = multiplicity 2 c"
    define c' where "c' = c div 2 ^ n"
    have "c = 2 ^n * c'"
      unfolding c'_def n_def by (simp add: multiplicity_dvd)
    have n_altdef: "n = multiplicity 2 c1"
      using <odd q> by (auto simp: n_def c multiplicity_prime_elem_times_other)
```

```
have "odd c'"
      unfolding c'_def n_def using <c > 0> multiplicity_decompose[of c
2] by auto
    define T where "T = (\lambda c. (\sum v = 1... < (a + 1) \text{ div } 2. \text{ [real\_of\_int (2)]})
* c * v) / real_of_int a|))"
    have "[2 * a * dedekind_sum' a c = a^2 + c^2 + 1 + 5 * c - 4 * c * T
c] \pmod{2 (n + 3)}"
      unfolding n_def T_def using <a > 0> <c > 0> <coprime a c> <even
c> <odd a>
      by (intro dedekind_sum'_cong_power_of_two') auto
    moreover have "[2 * a * dedekind_sum' a c1 * q = (a^2 + c1^2 + 1 + 1)
5 * c1 - 4 * c1 * T c1) * q
                      (mod 2 ^ (n + 3))"
      unfolding n_altdef T_def using <a > 0> <c1 > 0> <coprime a c1>
<even c1> <odd a>
      by (intro cong_mult[of _ _ q] dedekind_sum'_cong_power_of_two')
auto
    ultimately have "[r * a * \delta]" =
                         r * ((a^2 + c^2 + 1 + 5 * c - 4 * c * T c) - a *
(a + d) -
                         ((a^2 + c1^2 + 1 + 5 * c1 - 4 * c1 * T c1) * q -
a * (a + d) * q))]
                       (mod 2 ^ (n + 3))"
      unfolding r_a_\delta by (intro cong_mult[of r] cong_diff cong_refl)
(auto simp: mult_ac)
    also have "r * ((a^2 + c^2 + 1 + 5 * c - 4 * c * T c) - a * (a + d)
                         ((a^2 + c1^2 + 1 + 5 * c1 - 4 * c1 * T c1) * q -
a * (a + d) * q)) =
                r*(q-1) * (a * d - 1 + c * c1) - 4*c*r * (T c - T c1)"
      by (simp add: algebra_simps c power2_eq_square)
    also have "r * (q - 1) = 24"
      using q by (auto simp: r_def)
    also have "a * d - 1 = b * c"
      using det by (simp add: algebra_simps)
    also have "24 * (b * c + c * c1) = 24 * c * (b + c1)"
      by (simp add: algebra_simps)
    also have "[24 * c * (b + c1) - 4 * c * r * (T c - T c1) = 0 - 0]
(mod 2 ^ (n + 3))"
    proof (intro cong_diff)
      have "2 \hat{} (n + 3) dvd 2 \hat{} (n + 3) * (3 * c' * (b + c1))"
        using dvd\_triv\_left by blast
      also have "... = 24 * c * (b + c1)"
        by (simp add: \langle c = 2 \ \hat{} \ n * c' \rangle mult_ac power_add)
      finally show "[24 * c * (b + c1) = 0] (mod 2 ^ (n + 3))"
        by (simp add: cong_0_iff)
    \mathbf{next}
```

```
have "4 * 2 ^ n * 2 * 1 dvd 4 * c * r * (T c - T c1)"
         using \langle even \ r \rangle by (intro mult_dvd_mono) (auto simp: \langle c = 2 \ n \rangle
* c'>)
      thus "[4 * c * r * (T c - T c1) = 0] \pmod{2 (n + 3)}"
         by (simp add: power_add cong_0_iff mult_ac)
    finally have "2 \hat{} (n + 3) dvd r * \delta' * a"
      by (simp add: cong_0_iff mult_ac)
    hence "2 \hat{} (n + 3) dvd r * \delta'"
      using <odd a> by (subst (asm) coprime_dvd_mult_left_iff) auto
    have "2 \hat{} (n + 3) * (3 * c') dvd r * \delta'"
    proof (rule divides_mult)
      show "coprime (2 \hat{ } (n + 3)) (3 * c')"
         using <odd c'> by simp
      show "2 ^ (n + 3) dvd r * \delta'"
         by fact
      have "3 * c' dvd 3 * c"
         by (auto simp: \langle c = 2 \hat{n} * c' \rangle)
      also have "3 * c dvd r * \delta'"
         by fact
      finally show "3 * c' dvd r * \delta'".
    thus ?thesis
      by (simp add: <c = 2 ^ n * c'> power_add mult_ac)
  qed
qed
theorem dedekind_sum_diff_even_int:
  fixes a b c d :: int assumes det: "a * d - b * c = 1"
  fixes q c1 r :: int and \delta' :: "int \Rightarrow int" and \delta :: "int \Rightarrow real"
  assumes q: "q \in \{3, 5, 7, 13\}" and "c1 > 0"
  assumes c: "c = q * c1"
  defines "r \equiv 24 \text{ div } (q - 1)"
  defines "\delta' \equiv (\lambdaa. 2 * dedekind_sum' a c - (a + d) - (2 * q * dedekind_sum')
a c1 - (a + d) * q))"
  defines "\delta \equiv (\lambda a. \text{ dedekind\_sum a } c - (a+d)/(12*c) - (\text{dedekind\_sum a})
c1 - (a+d)/(12*c1))"
            "of_int (\delta' a) = 12 * c * \delta a"
  \mathbf{shows}
            "24 * c dvd r * \delta' a"
    and
    and
            "real_of_int r * \delta a / 2 \in \mathbb{Z}"
proof -
  define a' t where "a' = a mod c" and "t = a div c"
  have a'_eq: "a' = a - t * c"
    by (simp add: a'_def t_def algebra_simps)
  have cong1: "[a' = a] \pmod{c}"
    by (simp add: a'_def)
  hence cong2: "[a' = a] (mod c1)"
```

```
using c by (metis cong_modulus_mult mult.commute)
  have "q > 0"
    using q by auto
  have "c > 0"
    using q and \langle c1 \rangle 0> and c by auto
  have "[a * d - b * 0 = a * d - b * c] \pmod{c}"
    by (intro cong_diff cong_mult cong_refl) (auto simp: Cong.cong_def)
  also have "a * d - b * c = 1"
    by fact
  finally have "[a * d = 1] (mod c)"
    by simp
  hence "coprime a c"
    using coprime_iff_invertible_int by blast
  have "coprime a c1"
    using <coprime a c> c by auto
  have "a' \geq 0"
    using \langle c \rangle 0> by (auto simp: a'_def)
  moreover have "a' \neq 0"
    using <coprime a c> q by (auto simp: a'_def c)
  ultimately have "a' > 0"
    by linarith
  have 1: "dedekind_sum a' c = dedekind_sum a c" using cong1
    by (rule dedekind_sum_cong) (use <coprime a c > in <auto simp: a'_def
coprime_commute>)
  have 2: "dedekind_sum' a' c = dedekind_sum' a c" using cong1
    by (rule dedekind_sum'_cong) (use <coprime a c> in <auto simp: a'_def
coprime_commute>)
  have 3: "dedekind_sum a' c1 = dedekind_sum a c1" using cong2
    by (rule dedekind_sum_cong) (use <coprime a c1> in <auto simp: a'_def
coprime_commute>)
  have 4: "dedekind_sum' a' c1 = dedekind_sum' a c1" using cong2
    by (rule dedekind_sum'_cong) (use <coprime a c1> in <auto simp: a'_def
coprime commute>)
  have \delta_{eq}: "\delta a' = \delta a - t * (q - 1) / 12"
    unfolding \delta_{def} 1 3 using \langle c \rangle 0> \langle c1 \rangle 0> \langle q \rangle 0>
    by (simp add: field_simps c a'_eq)
  have \delta'_{eq}: "\delta'_{a} a' = \delta'_{a} a - c*t*(q-1)"
    unfolding \delta'_def 2 4 using \langle c \rangle 0 \rangle \langle c1 \rangle 0 \rangle \langle q \rangle 0 \rangle
    by (simp add: field_simps c a'_eq)
  have det': "a' * d - (b - t * d) * c = 1"
    using det by (simp add: a'_eq algebra_simps)
  have "of_int (\delta, a) = 12 * c * \delta a"
    unfolding \delta'_{def} \delta_{def} using det' \langle c1 \rangle 0 \rangle q c \langle a' \rangle 0 \rangle
```

```
by (intro dedekind_sum_diff_even_int_aux[of a' d "b - t * d" c]) auto
  thus of_int_\delta': "of_int (\delta' a) = 12 * c * \delta a"
    by (simp add: \delta_{eq} \delta'_{eq} field_simps)
  have "24 * c dvd r * \delta' a'"
    unfolding r_def \delta'_def using det' \langle c1 \rangle 0 \rangle q c \langle a' \rangle 0 \rangle
    by (intro dedekind_sum_diff_even_int_aux[of a' d "b - t * d" c]) auto
  also have "r * \delta' a' = r * \delta' a - (q - 1) * r * c * t"
    by (simp add: \delta '_eq algebra_simps)
  also have "(q - 1) * r = 24"
    unfolding r_{def} using q by auto
  also have "24 * c dvd r * \delta' a - 24 * c * t \longleftrightarrow 24 * c dvd r * \delta' a"
    by (rule dvd_diff_left_iff) auto
  finally show "24 * c dvd r * \delta' a" .
  then obtain m where m: "r * \delta, a = 24 * c * m"
    by auto
  hence "real_of_int (r * \delta' a) = real_of_int (24 * c * m)"
    by (simp only: )
  hence "real_of_int r * \delta a / 2 = of_int m"
    using \langle c \rangle 0> by (simp add: of_int_\delta' field_simps)
  also have "... \in \mathbb{Z}"
    by simp
  finally show "real_of_int r * \delta a / 2 \in \mathbb{Z}" .
qed
no_notation dedekind_frac ("\langle \rangle")
end
```

## References

[1] T. M. Apostol. Modular Functions and Dirichlet Series in Number Theory. Graduate Texts in Mathematics. Springer, 1990.