Decreasing-Diagrams-II

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Abstract

This theory formalizes a commutation version of decreasing diagrams for Church-Rosser modulo. The proof follows Felgenhauer and van Oostrom (RTA 2013). The theory also provides important specializations, in particular van Oostrom's conversion version (TCS 2008) of decreasing diagrams.

We follow the development described in [1]: Conversions are mapped to Greek strings, and we prove that whenever a local peak (or cliff) is replaced by a joining sequence from a locally decreasing diagram, then the corresponding Greek strings become smaller in a specially crafted well-founded order on Greek strings. Once there are no more local peaks or cliffs are left, the result is a valley that establishes the Church-Rosser modulo property.

As special cases we provide non-commutation versions and the conversion version of decreasing diagrams by van Oostrom [3]. We also formalize extended decreasingness [2].

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1 Preliminaries

theory *Decreasing-Diagrams-II-Aux*
imports
Well-Quasi-Orders.Multiset-Extension
Well-Quasi-Orders.Well-Quasi-Orders
begin

1.1 Trivialities

lemma *asymI2*: $(\bigwedge a\ b.\ (a,b) \in R \Longrightarrow (b,a) \notin R) \Longrightarrow asym\ R$
⟨*proof*⟩

abbreviation *strict-order R* ≡ *irrefl R* ∧ *trans R*

lemma *order-asym*: *trans R* ⟹ *asym R* = *irrefl R*
⟨*proof*⟩

lemma *strict-order-strict*: *strict-order q* ⟹ *strict* $(\lambda a\ b.\ (a,\ b) \in q^{=}) = (\lambda a\ b.\ (a,\ b) \in q)$
⟨*proof*⟩

lemma *mono-lex1*: *mono* (λ*r*. *lex-prod r s*)
⟨*proof*⟩

lemma *mono-lex2*: *mono* (*lex-prod r*)
⟨*proof*⟩

lemma *irrefl-lex-prod*: *irrefl R* ⟹ *irrefl S* ⟹ *irrefl* (*R* <∗*lex*∗> *S*)
⟨*proof*⟩

lemmas *converse-inward* = *rtrancl-converse*[*symmetric*] *converse-Un converse-UNION converse-relcomp*
converse-converse converse-Id

1.2 Complete lattices and least fixed points

context *complete-lattice*
begin

1.2.1 A chain-based induction principle

abbreviation *set-chain* :: *'a set* ⇒ *bool* **where**
set-chain C ≡ ∀ *x* ∈ *C*. ∀ *y* ∈ *C*. *x* ≤ *y* ∨ *y* ≤ *x*

lemma *lfp-chain-induct*:
assumes *mono*: *mono f*
and *step*: ⋀*x*. *P x* ⟹ *P* (*f x*)
and *chain*: ⋀*C*. *set-chain C* ⟹ ∀ *x* ∈ *C*. *P x* ⟹ *P* (*Sup C*)
shows *P* (*lfp f*)
⟨*proof*⟩

1.2.2 Preservation of transitivity, asymmetry, irreflexivity by suprema

lemma *trans-Sup-of-chain*:
assumes *set-chain C* **and** *trans*: ⋀*R*. *R* ∈ *C* ⟹ *trans R*
shows *trans* (*Sup C*)
⟨*proof*⟩

lemma *asym-Sup-of-chain*:
assumes *set-chain C* **and** *asym*: ⋀ *R*. *R* ∈ *C* ⟹ *asym R*
shows *asym* (*Sup C*)
⟨*proof*⟩

lemma *strict-order-lfp*:
assumes *mono f* **and** ⋀*R*. *strict-order R* ⟹ *strict-order* (*f R*)
shows *strict-order* (*lfp f*)
⟨*proof*⟩

lemma *trans-lfp*:
assumes *mono f* **and** ⋀*R*. *trans R* ⟹ *trans* (*f R*)
shows *trans* (*lfp f*)
⟨*proof*⟩

end

1.3 Multiset extension

lemma *mulex-iff-mult*: *mulex r M N* ⟷ (*M*,*N*) ∈ *mult* {(*M*,*N*) . *r M N*}
⟨*proof*⟩

lemma *multI*:
  assumes *trans r M = I + K N = I + J J ≠ {#} ∀k ∈ set-mset K. ∃j ∈ set-mset J. (k,j) ∈ r*
  shows *(M,N) ∈ mult r*
⟨*proof*⟩

lemma *multE*:
  assumes *trans r* **and** *(M,N) ∈ mult r*
  obtains *I J K* **where** *M = I + K N = I + J J ≠ {#} ∀k ∈ set-mset K. ∃j ∈ set-mset J. (k,j) ∈ r*
⟨*proof*⟩

lemma *mult-on-union*: *(M,N) ∈ mult r ⟹ (K + M, K + N) ∈ mult r*
⟨*proof*⟩

lemma *mult-on-union'*: *(M,N) ∈ mult r ⟹ (M + K, N + K) ∈ mult r*
⟨*proof*⟩

lemma *mult-on-add-mset*: *(M,N) ∈ mult r ⟹ (add-mset k M, add-mset k N) ∈ mult r*
⟨*proof*⟩

lemma *mult-empty*[*simp*]: *(M,{#}) ∉ mult R*
⟨*proof*⟩

lemma *mult-singleton*[*simp*]: *(x, y) ∈ r ⟹ (add-mset x M, add-mset y M) ∈ mult r*
⟨*proof*⟩

lemma *empty-mult*[*simp*]: *({#},N) ∈ mult R ⟷ N ≠ {#}*
⟨*proof*⟩

lemma *trans-mult*: *trans (mult R)*
⟨*proof*⟩

lemma *strict-order-mult*:
  assumes *irrefl R* **and** *trans R*
  shows *irrefl (mult R)* **and** *trans (mult R)*
⟨*proof*⟩

lemma *mult-of-image-mset*:
  assumes *trans R* **and** *trans R'*
  and *⋀x y. x ∈ set-mset N ⟹ y ∈ set-mset M ⟹ (x,y) ∈ R ⟹ (f x, f y) ∈ R'*
  and *(N, M) ∈ mult R*
  shows *(image-mset f N, image-mset f M) ∈ mult R'*
⟨*proof*⟩

1.4 Incrementality of *mult1* and *mult*

lemma *mono-mult1*: *mono mult1*
⟨*proof*⟩

lemma *mono-mult*: *mono mult*
⟨*proof*⟩

1.5 Well-orders and well-quasi-orders

lemma *wf-iff-wfp-on*:
wf p ⟷ *wfp-on* (λ*a b*. (*a*, *b*) ∈ *p*) *UNIV*
⟨*proof*⟩

lemma *well-order-implies-wqo*:
assumes *well-order r*
shows *wqo-on* (λ*a b*. (*a*, *b*) ∈ *r*) *UNIV*
⟨*proof*⟩

1.6 Splitting lists into prefix, element, and suffix

fun *list-splits* :: *'a list* ⇒ (*'a list* × *'a* × *'a list*) *list* **where**
list-splits [] = []
| *list-splits* (*x* # *xs*) = ([], *x*, *xs*) # *map* (λ(*xs*, *x'*, *xs'*). (*x* # *xs*, *x'*, *xs'*)) (*list-splits xs*)

lemma *list-splits-empty*[*simp*]:
list-splits xs = [] ⟷ *xs* = []
⟨*proof*⟩

lemma *elem-list-splits-append*:
assumes (*ys*, *y*, *zs*) ∈ *set* (*list-splits xs*)
shows *ys* @ [*y*] @ *zs* = *xs*
⟨*proof*⟩

lemma *elem-list-splits-length*:
assumes (*ys*, *y*, *zs*) ∈ *set* (*list-splits xs*)
shows *length ys* < *length xs* **and** *length zs* < *length xs*
⟨*proof*⟩

lemma *elem-list-splits-elem*:
assumes (*xs*, *y*, *ys*) ∈ *set* (*list-splits zs*)
shows *y* ∈ *set zs*
⟨*proof*⟩

lemma *list-splits-append*:
list-splits (*xs* @ *ys*) = *map* (λ(*xs'*, *x'*, *ys'*). (*xs'*, *x'*, *ys'* @ *ys*)) (*list-splits xs*) @
map (λ(*xs'*, *x'*, *ys'*). (*xs* @ *xs'*, *x'*, *ys'*)) (*list-splits ys*)
⟨*proof*⟩

lemma *list-splits-rev*:
list-splits (*rev xs*) = *map* (λ(*xs*, *x*, *ys*). (*rev ys*, *x*, *rev xs*)) (*rev* (*list-splits xs*))
⟨*proof*⟩

lemma *list-splits-map*:
list-splits (*map f xs*) = *map* (λ(*xs*, *x*, *ys*). (*map f xs*, *f x*, *map f ys*)) (*list-splits xs*)
⟨*proof*⟩

end

2 Decreasing Diagrams

theory *Decreasing-Diagrams-II*
imports
Decreasing-Diagrams-II-Aux
HOL−Cardinals.Wellorder-Extension
Abstract−Rewriting.Abstract-Rewriting
begin

2.1 Greek accents

datatype *accent* = *Acute* | *Grave* | *Macron*

lemma *UNIV-accent*: *UNIV* = { *Acute*, *Grave*, *Macron* }
⟨*proof*⟩

lemma *finite-accent*: *finite* (*UNIV* :: *accent set*)
⟨*proof*⟩

type-synonym *'a letter* = *accent* × *'a*

definition *letter-less* :: (*'a* × *'a*) *set* ⇒ (*'a letter* × *'a letter*) *set* **where**
[*simp*]: *letter-less R* = {(*a*,*b*). (*snd a*, *snd b*) ∈ *R*}

lemma *mono-letter-less*: *mono letter-less*
⟨*proof*⟩

2.2 Comparing Greek strings

type-synonym *'a greek* = *'a letter list*

definition *adj-msog* :: *'a greek* ⇒ *'a greek* ⇒ (*'a letter* × *'a greek*) ⇒ (*'a letter* × *'a greek*)
where
adj-msog xs zs l ≡
case l of (*y*,*ys*) ⇒ (*y*, *case fst y of Acute* ⇒ *ys* @ *zs* | *Grave* ⇒ *xs* @ *ys* | *Macron* ⇒ *ys*)

definition *ms-of-greek* :: *'a greek* ⇒ (*'a letter* × *'a greek*) *multiset* **where**
ms-of-greek as = *mset*
(*map* (λ(*xs*, *y*, *zs*) ⇒ *adj-msog xs zs* (*y*, [])) (*list-splits as*))

lemma *adj-msog-adj-msog*[*simp*]:
adj-msog xs zs (*adj-msog xs' zs' y*) = *adj-msog* (*xs* @ *xs'*) (*zs'* @ *zs*) *y*
⟨*proof*⟩

lemma *compose-adj-msog*[*simp*]: *adj-msog xs zs* ∘ *adj-msog xs' zs'* = *adj-msog* (*xs* @ *xs'*) (*zs'* @ *zs*)
⟨*proof*⟩

lemma *adj-msog-single*:
adj-msog xs zs (*x*,[]) = (*x*, (*case fst x of Grave* ⇒ *xs* | *Acute* ⇒ *zs* | *Macron* ⇒ []))
⟨*proof*⟩

lemma *ms-of-greek-elem*:
assumes (*x*,*xs*) ∈ *set-mset* (*ms-of-greek ys*)
shows *x* ∈ *set ys*
⟨*proof*⟩

lemma *ms-of-greek-shorter*:
assumes (*x*, *t*) ∈# *ms-of-greek s*
shows *length s* > *length t*
⟨*proof*⟩

lemma *msog-append*: *ms-of-greek* (*xs* @ *ys*) = *image-mset* (*adj-msog* [] *ys*) (*ms-of-greek xs*) +
image-mset (*adj-msog xs* []) (*ms-of-greek ys*)
⟨*proof*⟩

definition *nest* :: (*'a* × *'a*) *set* ⇒ (*'a greek* × *'a greek*) *set* ⇒ (*'a greek* × *'a greek*) *set* **where**
[*simp*]: *nest r s* = {(*a*,*b*). (*ms-of-greek a*, *ms-of-greek b*) ∈ *mult* (*letter-less r* <∗*lex*∗> *s*)}

lemma *mono-nest*: *mono* (*nest r*)
⟨*proof*⟩

lemma *nest-mono*[*mono-set*]: *x* ⊆ *y* ⟹ (*a*,*b*) ∈ *nest r x* ⟶ (*a*,*b*) ∈ *nest r y*
⟨*proof*⟩

definition *greek-less* :: (*'a* × *'a*) *set* ⇒ (*'a greek* × *'a greek*) *set* **where**
greek-less r = *lfp* (*nest r*)

lemma *greek-less-unfold*:
greek-less r = *nest r* (*greek-less r*)
⟨*proof*⟩

2.3 Preservation of strict partial orders

lemma *strict-order-letter-less*:
assumes *strict-order r*
shows *strict-order* (*letter-less r*)
⟨*proof*⟩

lemma *strict-order-nest*:
assumes *r*: *strict-order r* **and** *R*: *strict-order R*
shows *strict-order* (*nest r R*)
⟨*proof*⟩

lemma *strict-order-greek-less*:
assumes *strict-order r*
shows *strict-order* (*greek-less r*)
⟨*proof*⟩

lemma *trans-letter-less*:
assumes *trans r*
shows *trans* (*letter-less r*)
⟨*proof*⟩

lemma *trans-order-nest*: *trans* (*nest r R*)
⟨*proof*⟩

lemma *trans-greek-less*[*simp*]: *trans* (*greek-less r*)
⟨*proof*⟩

lemma *mono-greek-less*: *mono greek-less*
⟨*proof*⟩

2.4 Involution

definition *inv-letter* :: *'a letter* ⇒ *'a letter* **where**
inv-letter l ≡
case l of (*a*, *x*) ⇒ (*case a of Grave* ⇒ *Acute* | *Acute* ⇒ *Grave* | *Macron* ⇒ *Macron*, *x*)

lemma *inv-letter-pair*[*simp*]:
inv-letter (*a*, *x*) = (*case a of Grave* ⇒ *Acute* | *Acute* ⇒ *Grave* | *Macron* ⇒ *Macron*, *x*)
⟨*proof*⟩

lemma *snd-inv-letter*[*simp*]:
snd (*inv-letter x*) = *snd x*
⟨*proof*⟩

lemma *inv-letter-invol*[*simp*]:
inv-letter (*inv-letter x*) = *x*
⟨*proof*⟩

lemma *inv-letter-mono*[*simp*]:
assumes $(x, y) \in$ *letter-less r*
shows $(\textit{inv-letter } x, \textit{inv-letter } y) \in \textit{letter-less } r$
⟨*proof*⟩

definition *inv-greek* :: *'a greek* ⇒ *'a greek* **where**
inv-greek s = *rev* (*map inv-letter s*)

lemma *inv-greek-invol*[*simp*]:
inv-greek (*inv-greek s*) = *s*
⟨*proof*⟩

lemma *inv-greek-append*:
inv-greek (*s* @ *t*) = *inv-greek t* @ *inv-greek s*
⟨*proof*⟩

definition *inv-msog* :: (*'a letter* × *'a greek*) *multiset* ⇒ (*'a letter* × *'a greek*) *multiset* **where**
inv-msog M = *image-mset* (λ(*x*, *t*). (*inv-letter x*, *inv-greek t*)) *M*

lemma *inv-msog-invol*[*simp*]:
inv-msog (*inv-msog M*) = *M*
⟨*proof*⟩

lemma *ms-of-greek-inv-greek*:
ms-of-greek (*inv-greek M*) = *inv-msog* (*ms-of-greek M*)
⟨*proof*⟩

lemma *inv-greek-mono*:
assumes *trans r* **and** (*s*, *t*) ∈ *greek-less r*
shows (*inv-greek s*, *inv-greek t*) ∈ *greek-less r*
⟨*proof*⟩

2.5 Monotonicity of *greek-less r*

lemma *greek-less-rempty*[*simp*]:
(*a*,[]) ∈ *greek-less r* ⟷ *False*
⟨*proof*⟩

lemma *greek-less-nonempty*:
assumes *b* ≠ []
shows (*a*,*b*) ∈ *greek-less r* ⟷ (*a*,*b*) ∈ *nest r* (*greek-less r*)
⟨*proof*⟩

lemma *greek-less-lempty*[*simp*]:
([],*b*) ∈ *greek-less r* ⟷ *b* ≠ []
⟨*proof*⟩

lemma *greek-less-singleton*:
(*a*, *b*) ∈ *letter-less r* ⟹ ([*a*], [*b*]) ∈ *greek-less r*
⟨*proof*⟩

lemma *ms-of-greek-cons*:
ms-of-greek (*x* # *s*) = {# *adj-msog* [] *s* (*x*,[]) #} + *image-mset* (*adj-msog* [*x*] []) (*ms-of-greek s*)
⟨*proof*⟩

lemma *greek-less-cons-mono*:
assumes *trans r*
shows (*s*, *t*) ∈ *greek-less r* ⟹ (*x* # *s*, *x* # *t*) ∈ *greek-less r*
⟨*proof*⟩

lemma *greek-less-app-mono2*:
assumes *trans r* **and** (*s*, *t*) ∈ *greek-less r*
shows (*p* @ *s*, *p* @ *t*) ∈ *greek-less r*
⟨*proof*⟩

lemma *greek-less-app-mono1*:
assumes *trans r* **and** (*s*, *t*) ∈ *greek-less r*
shows (*s* @ *p*, *t* @ *p*) ∈ *greek-less r*
⟨*proof*⟩

2.6 Well-founded-ness of *greek-less r*

lemma *greek-embed*:
assumes *trans r*
shows *list-emb* (λ*a b*. (*a*, *b*): *reflcl* (*letter-less r*)) *a b* ⟹ (*a*, *b*) ∈ *reflcl* (*greek-less r*)
⟨*proof*⟩

lemma *wqo-letter-less*:
assumes *t*: *trans r* **and** *w*: *wqo-on* (λ*a b*. (*a*, *b*) ∈ $r^=$) *UNIV*
shows *wqo-on* (λ*a b*. (*a*, *b*) ∈ (*letter-less r*)$^=$) *UNIV*
⟨*proof*⟩

lemma *wf-greek-less*:
assumes *wf r* **and** *trans r*
shows *wf* (*greek-less r*)
⟨*proof*⟩

2.7 Basic Comparisons

lemma *pairwise-imp-mult*:
assumes *N* ≠ {#} **and** ∀ *x* ∈ *set-mset M*. ∃ *y* ∈ *set-mset N*. (*x*, *y*) ∈ *r*
shows (*M*, *N*) ∈ *mult r*
⟨*proof*⟩

lemma *singleton-greek-less*:

assumes *as*: *snd ‘ set as* ⊆ *under r b*
shows (*as*, [(*a*,*b*)]) ∈ *greek-less r*
⟨*proof*⟩

lemma *peak-greek-less*:
assumes *as*: *snd ‘ set as* ⊆ *under r a* **and** *b′*: *b′* ∈ {[(*Grave*,*b*)],[]}
and *cs*: *snd ‘ set cs* ⊆ *under r a* ∪ *under r b* **and** *a′*: *a′* ∈ {[(*Acute*,*a*)],[]}
and *bs*: *snd ‘ set bs* ⊆ *under r b*
shows (*as* @ *b′* @ *cs* @ *a′* @ *bs*, [(*Acute*,*a*),(*Grave*,*b*)]) ∈ *greek-less r*
⟨*proof*⟩

lemma *rcliff-greek-less1*:
assumes *trans r*
and *as*: *snd ‘ set as* ⊆ *under r a* ∩ *under r b* **and** *b′*: *b′* ∈ {[(*Grave*,*b*)],[]}
and *cs*: *snd ‘ set cs* ⊆ *under r b* **and** *a′*: *a′* = [(*Macron*,*a*)]
and *bs*: *snd ‘ set bs* ⊆ *under r b*
shows (*as* @ *b′* @ *cs* @ *a′* @ *bs*, [(*Macron*,*a*),(*Grave*,*b*)]) ∈ *greek-less r*
⟨*proof*⟩

lemma *rcliff-greek-less2*:
assumes *trans r*
and *as*: *snd ‘ set as* ⊆ *under r a* **and** *b′*: *b′* ∈ {[(*Grave*,*b*)],[]}
and *cs*: *snd ‘ set cs* ⊆ *under r a* ∪ *under r b*
shows (*as* @ *b′* @ *cs*, [(*Macron*,*a*),(*Grave*,*b*)]) ∈ *greek-less r*
⟨*proof*⟩

lemma *snd-inv-greek* [*simp*]: *snd ‘ set* (*inv-greek as*) = *snd ‘ set as*
⟨*proof*⟩

lemma *lcliff-greek-less1*:
assumes *trans r*
and *as*: *snd ‘ set as* ⊆ *under r a* **and** *b′*: *b′* = [(*Macron*,*b*)]
and *cs*: *snd ‘ set cs* ⊆ *under r a* **and** *a′*: *a′* ∈ {[(*Acute*,*a*)],[]}
and *bs*: *snd ‘ set bs* ⊆ *under r a* ∩ *under r b*
shows (*as* @ *b′* @ *cs* @ *a′* @ *bs*, [(*Acute*,*a*),(*Macron*,*b*)]) ∈ *greek-less r*
⟨*proof*⟩

lemma *lcliff-greek-less2*:
assumes *trans r*
and *cs*: *snd ‘ set cs* ⊆ *under r a* ∪ *under r b* **and** *a′*: *a′* ∈ {[(*Acute*,*a*)],[]}
and *bs*: *snd ‘ set bs* ⊆ *under r b*
shows (*cs* @ *a′* @ *bs*, [(*Acute*,*a*),(*Macron*,*b*)]) ∈ *greek-less r*
⟨*proof*⟩

2.8 Labeled abstract rewriting

context
fixes *L R E* :: *′b* ⇒ *′a rel*
begin

definition *lstep* :: *'b letter* ⇒ *'a rel* **where**
  [*simp*]: *lstep x* = (*case x of* (*a*, *i*) ⇒ (*case a of Acute* ⇒ (*L i*)$^{-1}$ | *Grave* ⇒ *R i*
| *Macron* ⇒ *E i*))

fun *lconv* :: *'b greek* ⇒ *'a rel* **where**
  lconv [] = *Id*
| *lconv* (*x* # *xs*) = *lstep x O lconv xs*

lemma *lconv-append*[*simp*]:
  lconv (*xs* @ *ys*) = *lconv xs O lconv ys*
⟨*proof*⟩

lemma *conversion-join-or-peak-or-cliff*:
  obtains (*join*) *as bs cs* **where** *set as* ⊆ {*Grave*} **and** *set bs* ⊆ {*Macron*} **and**
set cs ⊆ {*Acute*}
    and *ds* = *as* @ *bs* @ *cs*
  | (*peak*) *as bs* **where** *ds* = *as* @ ([*Acute*] @ [*Grave*]) @ *bs*
  | (*lcliff*) *as bs* **where** *ds* = *as* @ ([*Acute*] @ [*Macron*]) @ *bs*
  | (*rcliff*) *as bs* **where** *ds* = *as* @ ([*Macron*] @ [*Grave*]) @ *bs*
⟨*proof*⟩

lemma *map-eq-append-split*:
  assumes *map f xs* = *ys1* @ *ys2*
  obtains *xs1 xs2* **where** *ys1* = *map f xs1 ys2* = *map f xs2 xs* = *xs1* @ *xs2*
⟨*proof*⟩

lemmas *map-eq-append-splits* = *map-eq-append-split map-eq-append-split*[*OF sym*]

abbreviation *conversion'* *M* ≡ ((⋃ *i* ∈ *M*. *R i*) ∪ (⋃ *i* ∈ *M*. *E i*) ∪ (⋃ *i* ∈ *M*. *L*
i)$^{-1}$)*
abbreviation *valley'* *M* ≡ (⋃ *i* ∈ *M*. *R i*)* *O* (⋃ *i* ∈ *M*. *E i*)* *O* ((⋃ *i* ∈ *M*. *L*
i)$^{-1}$)*

lemma *conversion-to-lconv*:
  assumes (*u*, *v*) ∈ *conversion' M*
  obtains *xs* **where** *snd* ‘ *set xs* ⊆ *M* **and** (*u*, *v*) ∈ *lconv xs*
⟨*proof*⟩

definition *lpeak* :: *'b rel* ⇒ *'b* ⇒ *'b* ⇒ *'b greek* ⇒ *bool* **where**
  lpeak r a b xs ⟷ (∃ *as b' cs a' bs*. *snd* ‘ *set as* ⊆ *under r a* ∧ *b'* ∈ {[(*Grave*,*b*)],[]}
∧
    snd ‘ *set cs* ⊆ *under r a* ∪ *under r b* ∧ *a'* ∈ {[(*Acute*,*a*)],[]} ∧
    snd ‘ *set bs* ⊆ *under r b* ∧ *xs* = *as* @ *b'* @ *cs* @ *a'* @ *bs*)

definition *lcliff* :: *'b rel* ⇒ *'b* ⇒ *'b* ⇒ *'b greek* ⇒ *bool* **where**
  lcliff r a b xs ⟷ (∃ *as b' cs a' bs*. *snd* ‘ *set as* ⊆ *under r a* ∧ *b'* = [(*Macron*,*b*)]
∧
    snd ‘ *set cs* ⊆ *under r a* ∧ *a'* ∈ {[(*Acute*,*a*)],[]} ∧

snd ‘ set bs ⊆ *under r a* ∩ *under r b* ∧ *xs* = *as* @ *b′* @ *cs* @ *a′* @ *bs*) ∨
(∃ *cs a′ bs. snd ‘ set cs* ⊆ *under r a* ∪ *under r b* ∧ *a′* ∈ {[(*Acute*,*a*)],[]} ∧
snd ‘ set bs ⊆ *under r b* ∧ *xs* = *cs* @ *a′* @ *bs*)

definition *rcliff* :: *′b rel* ⇒ *′b* ⇒ *′b* ⇒ *′b greek* ⇒ *bool* **where**
rcliff r a b xs ⟷ (∃ *as b′ cs a′ bs. snd ‘ set as* ⊆ *under r a* ∩ *under r b* ∧ *b′* ∈ {[(*Grave*,*b*)],[]} ∧
snd ‘ set cs ⊆ *under r b* ∧ *a′* = [(*Macron*,*a*)] ∧
snd ‘ set bs ⊆ *under r b* ∧ *xs* = *as* @ *b′* @ *cs* @ *a′* @ *bs*) ∨
(∃ *as b′ cs. snd ‘ set as* ⊆ *under r a* ∧ *b′* ∈ {[(*Grave*,*b*)],[]} ∧
snd ‘ set cs ⊆ *under r a* ∪ *under r b* ∧ *xs* = *as* @ *b′* @ *cs*)

lemma *dd-commute-modulo-conv*[*case-names wf trans peak lcliff rcliff*]:
assumes *wf r* **and** *trans r*
and *pk*: ⋀*a b s t u.* (*s*, *t*) ∈ *L a* ⟹ (*s*, *u*) ∈ *R b* ⟹ ∃ *xs. lpeak r a b xs* ∧ (*t*, *u*) ∈ *lconv xs*
and *lc*: ⋀*a b s t u.* (*s*, *t*) ∈ *L a* ⟹ (*s*, *u*) ∈ *E b* ⟹ ∃ *xs. lcliff r a b xs* ∧ (*t*, *u*) ∈ *lconv xs*
and *rc*: ⋀*a b s t u.* (*s*, *t*) ∈ (*E a*)$^{-1}$ ⟹ (*s*, *u*) ∈ *R b* ⟹ ∃ *xs. rcliff r a b xs* ∧ (*t*, *u*) ∈ *lconv xs*
shows *conversion′ UNIV* ⊆ *valley′ UNIV*
⟨*proof*⟩

3 Results

3.1 Church-Rosser modulo

Decreasing diagrams for Church-Rosser modulo, commutation version.

lemma *dd-commute-modulo*[*case-names wf trans peak lcliff rcliff*]:
assumes *wf r* **and** *trans r*
and *pk*: ⋀*a b s t u.* (*s*, *t*) ∈ *L a* ⟹ (*s*, *u*) ∈ *R b* ⟹
(*t*, *u*) ∈ *conversion′* (*under r a*) *O* (*R b*)$^{=}$ *O conversion′* (*under r a* ∪ *under r b*) *O*
((*L a*)$^{-1}$)$^{=}$ *O conversion′* (*under r b*)
and *lc*: ⋀*a b s t u.* (*s*, *t*) ∈ *L a* ⟹ (*s*, *u*) ∈ *E b* ⟹
(*t*, *u*) ∈ *conversion′* (*under r a*) *O E b O conversion′* (*under r a*) *O*
((*L a*)$^{-1}$)$^{=}$ *O conversion′* (*under r a* ∩ *under r b*) ∨
(*t*, *u*) ∈ *conversion′* (*under r a* ∪ *under r b*) *O* ((*L a*)$^{-1}$)$^{=}$ *O conversion′* (*under r b*)
and *rc*: ⋀*a b s t u.* (*s*, *t*) ∈ (*E a*)$^{-1}$ ⟹ (*s*, *u*) ∈ *R b* ⟹
(*t*, *u*) ∈ *conversion′* (*under r a* ∩ *under r b*) *O* (*R b*)$^{=}$ *O conversion′* (*under r b*) *O*
E a O conversion′ (*under r b*) ∨
(*t*, *u*) ∈ *conversion′* (*under r a*) *O* (*R b*)$^{=}$ *O conversion′* (*under r a* ∪ *under r b*)
shows *conversion′ UNIV* ⊆ *valley′ UNIV*
⟨*proof*⟩

end

Decreasing diagrams for Church-Rosser modulo.

lemma *dd-cr-modulo*[*case-names wf trans symE peak cliff*]:
assumes *wf r* **and** *trans r* **and** *E*: $\bigwedge i.\ sym\ (E\ i)$
and *pk*: $\bigwedge a\ b\ s\ t\ u.\ (s,\ t) \in L\ a \Longrightarrow (s,\ u) \in L\ b \Longrightarrow$
$(t,\ u) \in conversion'\ L\ L\ E\ (under\ r\ a)\ O\ (L\ b)^{=}\ O\ conversion'\ L\ L\ E\ (under\ r\ a \cup under\ r\ b)\ O$
$((L\ a)^{-1})^{=}\ O\ conversion'\ L\ L\ E\ (under\ r\ b)$
and *cl*: $\bigwedge a\ b\ s\ t\ u.\ (s,\ t) \in L\ a \Longrightarrow (s,\ u) \in E\ b \Longrightarrow$
$(t,\ u) \in conversion'\ L\ L\ E\ (under\ r\ a)\ O\ E\ b\ O\ conversion'\ L\ L\ E\ (under\ r\ a)\ O$
$((L\ a)^{-1})^{=}\ O\ conversion'\ L\ L\ E\ (under\ r\ a \cap under\ r\ b) \vee$
$(t,\ u) \in conversion'\ L\ L\ E\ (under\ r\ a \cup under\ r\ b)\ O\ ((L\ a\ )^{-1})^{=}\ O\ conversion'\ L\ L\ E\ (under\ r\ b)$
shows $conversion'\ L\ L\ E\ UNIV \subseteq valley'\ L\ L\ E\ UNIV$
⟨*proof*⟩

3.2 Commutation and confluence

abbreviation $conversion''\ L\ R\ M \equiv ((\bigcup i \in M.\ R\ i) \cup (\bigcup i \in M.\ L\ i)^{-1})^{*}$
abbreviation $valley''\ L\ R\ M \equiv (\bigcup i \in M.\ R\ i)^{*}\ O\ ((\bigcup i \in M.\ L\ i)^{-1})^{*}$

Decreasing diagrams for commutation.

lemma *dd-commute*[*case-names wf trans peak*]:
assumes *wf r* **and** *trans r*
and *pk*: $\bigwedge a\ b\ s\ t\ u.\ (s,\ t) \in L\ a \Longrightarrow (s,\ u) \in R\ b \Longrightarrow$
$(t,\ u) \in conversion''\ L\ R\ (under\ r\ a)\ O\ (R\ b)^{=}\ O\ conversion''\ L\ R\ (under\ r\ a \cup under\ r\ b)\ O$
$((L\ a)^{-1})^{=}\ O\ conversion''\ L\ R\ (under\ r\ b)$
shows $commute\ (\bigcup i.\ L\ i)\ (\bigcup i.\ R\ i)$
⟨*proof*⟩

Decreasing diagrams for confluence.

lemmas *dd-cr*[*case-names wf trans peak*] =
dd-commute[*of - L L* **for** *L*, *unfolded CR-iff-self-commute*[*symmetric*]]

3.3 Extended decreasing diagrams

context
fixes *r q* :: *'b rel*
assumes *wf r* **and** *trans r* **and** *trans q* **and** *refl q* **and** *compat*: $r\ O\ q \subseteq r$
begin

private abbreviation (*input*) *down* :: $('b \Rightarrow 'a\ rel) \Rightarrow ('b \Rightarrow 'a\ rel)$ **where**
$down\ L \equiv \lambda i.\ \bigcup j \in under\ q\ i.\ L\ j$

private lemma *Union-down*: $(\bigcup i.\ down\ L\ i) = (\bigcup i.\ L\ i)$
⟨*proof*⟩

Extended decreasing diagrams for commutation.

lemma *edd-commute*[*case-names wf transr transq reflq compat peak*]:
assumes *pk*: $\bigwedge a\ b\ s\ t\ u.\ (s,\ t) \in L\ a \Longrightarrow (s,\ u) \in R\ b \Longrightarrow$
$(t,\ u) \in \mathit{conversion}''\ L\ R\ (\mathit{under}\ r\ a)\ O\ (\mathit{down}\ R\ b)^{=}\ O\ \mathit{conversion}''\ L\ R\ (\mathit{under}\ r\ a \cup \mathit{under}\ r\ b)\ O$
$((\mathit{down}\ L\ a)^{-1})^{=}\ O\ \mathit{conversion}''\ L\ R\ (\mathit{under}\ r\ b)$
shows *commute* $(\bigcup i.\ L\ i)\ (\bigcup i.\ R\ i)$
⟨*proof*⟩

Extended decreasing diagrams for confluence.

lemmas *edd-cr*[*case-names wf transr transq reflq compat peak*] =
edd-commute[*of L L* **for** *L*, *unfolded CR-iff-self-commute*[*symmetric*]]

end

end

References

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