

Coproduct Measure

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Abstract

This entry formalizes the coproduct measure. Let I be a set and $\{M_i\}_{i \in I}$ measurable spaces. The σ -algebra on $\coprod_{i \in I} M_i = \{(i, x) \mid i \in I \wedge x \in M_i\}$ is defined as the least one making $(\lambda x. (i, x))$ measurable for all $i \in I$. Let μ_i be measures on M_i for all $i \in I$ and A a measurable set of $\coprod_{i \in I} M_i$. The coproduct measure $\coprod_{i \in I} \mu_i$ is defined as follows:

$$\left(\coprod_{i \in I} \mu_i\right)(A) = \sum_{i \in I} \mu_i(A_i), \quad \text{where } A_i = \{x \mid (i, x) \in A\}.$$

We also prove the relationship with coproduct quasi-Borel spaces: the functor $R : \mathbf{Meas} \rightarrow \mathbf{QBS}$ preserves countable coproducts.

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1 Preliminaries

```

theory Lemmas-Coproduct-Measure
  imports HOL-Probability.Probability
           Standard-Borel-Spaces.Abstract-Metrizable-Topology
begin

lemma metrizable-space-metric-space:
  assumes  $d$ :Metric-space UNIV  $d$  Metric-space.mtopology UNIV  $d$  = euclidean
  shows class.metric-space  $d$  ( $\prod e \in \{0 < ..\}$ . principal  $\{(x, y). d\ x\ y < e\}$ ) open
proof -
  interpret Metric-space UNIV  $d$  by fact
  show ?thesis
proof
  show open  $U \iff (\forall x \in U. \forall_F (x', y) \text{ in } \prod e \in \{0 < ..\}. \text{principal } \{(F, y). d\ F\ y < e\}. x' = x \implies y \in U)$  for  $U$ 
  proof(subst eventually-INF-base)
    show  $a \in \{0 < ..\} \implies b \in \{0 < ..\} \implies \exists x \in \{0 < ..\}. \text{principal } \{(F, y). d\ F\ y < a\} \leq \text{principal } \{(F, y). d\ F\ y < b\} \cap \text{principal } \{(F, y). d\ F\ y < b\}$  for  $a\ b$ 
    by(auto intro!: be_xI[where  $x = \min\ a\ b$ ])
  next
    show open  $U \iff (\forall x \in U. \exists b \in \{0 < ..\}. \forall_F (x', y) \text{ in } \text{principal } \{(F, y). d\ F\ y < b\}. x' = x \implies y \in U)$ 
    by(fastforce simp: open_in_mtopology[simplified  $d(2)$ ,simplified] eventually-principal)
  qed simp
qed(auto simp: triangle')
qed

corollary metrizable-space-metric-space-ex:
  assumes metrizable-space (euclidean :: 'a :: topological-space topology)
  shows  $\exists (d :: 'a \Rightarrow 'a \Rightarrow \text{real})\ F. \text{class.metric-space } d\ F\ \text{open}$ 
proof -
  from assms obtain  $d :: 'a \Rightarrow 'a \Rightarrow \text{real}$  where Metric-space UNIV  $d$  Metric-space.mtopology UNIV  $d$  = euclidean
  by (metis Metric-space.topspace-mtopology metrizable-space-def topspace-euclidean)
from metrizable-space-metric-space[OF this] show ?thesis

```

by blast
qed

lemma *completely-metrizable-space-metric-space:*

assumes *Metric-space* (*UNIV* :: 'a :: *topological-space set*) *d Metric-space.mtopology UNIV d = euclidean Metric-space.mcomplete UNIV d*

shows *class.complete-space d* ($\prod e \in \{0 < ..\}$. *principal* $\{(x,y). d\ x\ y < e\}$) *open*

proof –

interpret *Metric-space UNIV d by fact*

interpret *m:metric-space d* $\prod e \in \{0 < ..\}$. *principal* $\{(x,y). d\ x\ y < e\}$ *open*

by(*auto intro!*: *metrizable-space-metric-space assms*)

have [*simp*]:*topological-space.convergent* (*open* :: 'a *set* \Rightarrow *bool*) = *convergent*

proof

fix *x* :: *nat* \Rightarrow 'a

have *:*class.topological-space* (*open* :: 'a *set* \Rightarrow *bool*)

by *standard auto*

show *topological-space.convergent open x = convergent x*

by(*simp add*: *topological-space.convergent-def*[*OF* *] *topological-space.nhds-def*[*OF* *] *convergent-def nhds-def*)

qed

show *?thesis*

apply *unfold-locales*

using *assms*(3) by(*auto simp*: *mcomplete-def assms*(2) *MCauchy-def m.Cauchy-def convergent-def*)

qed

lemma *completely-metrizable-space-metric-space-ex:*

assumes *completely-metrizable-space* (*euclidean* :: 'a :: *topological-space topology*)

shows $\exists (d :: 'a \Rightarrow 'a \Rightarrow \text{real}) F.$ *class.complete-space d F open*

proof –

from *assms* **obtain** *d* :: 'a \Rightarrow 'a \Rightarrow *real* **where** *Metric-space UNIV d Metric-space.mtopology UNIV d = euclidean Metric-space.mcomplete UNIV d*

by (*metis Metric-space.topspace-mtopology completely-metrizable-space-def topspace-euclidean*)

from *completely-metrizable-space-metric-space*[*OF this*] **show** *?thesis*

by *blast*

qed

1.1 Polishness of Extended Reals and Non-Negative Extended Reals

We instantiate *polish-space* for *ereal* and *ennreal* with *non-canonical* metrics in order to change the order of *infsum* using the lemma *infsum-Sigma*.

instantiation *ereal* :: *metric-space*

begin

definition *dist-ereal* :: *ereal* \Rightarrow *ereal* \Rightarrow *real*

where *dist-ereal* \equiv *SOME d. Metric-space UNIV d \wedge*

Metric-space.mtopology UNIV d = euclidean \wedge
Metric-space.mcomplete UNIV d

definition *uniformity-ereal* :: (*ereal* \times *ereal*) filter
where *uniformity-ereal* $\equiv \prod_{e \in \{0 < ..\}}$. principal $\{(x,y). \text{dist } x \ y < e\}$

instance

proof –

let *?open* = *open* :: *ereal set* \Rightarrow *bool*

interpret *c:complete-space dist uniformity ?open*

proof –

have $\exists d.$ *Metric-space (UNIV :: ereal set) d* \wedge
Metric-space.mtopology UNIV d = euclidean \wedge
Metric-space.mcomplete UNIV d

by (*metis Polish-space-ereal Metric-space.topspace-mtopology Polish-space-def completely-metrizable-space-def topspace-euclidean*)

hence *Metric-space (UNIV :: ereal set) dist* \wedge
Metric-space.mtopology (UNIV :: ereal set) dist = euclidean \wedge
Metric-space.mcomplete (UNIV :: ereal set) dist

unfolding *dist-ereal-def* **by**(*rule someI-ex*)

with *completely-metrizable-space-metric-space* **show** *class.complete-space dist uniformity ?open*

by(*fastforce simp: uniformity-ereal-def*)

qed

have [*simp*]:*topological-space.convergent ?open = convergent*

proof

fix *x* :: *nat* \Rightarrow *ereal*

have **:class.topological-space ?open*

by *standard auto*

show *topological-space.convergent open x = convergent x*

by(*simp add: topological-space.convergent-def topological-space.nhds-def * convergent-def nhds-def*)

qed

show *OFCLASS(ereal, metric-space-class)*

by *standard (use uniformity-ereal-def c.open-uniformity c.dist-triangle2 c.Cauchy-convergent*

in auto)

qed

end

instantiation *ereal* :: *polish-space*

begin

instance

proof

let *?open* = *open* :: *ereal set* \Rightarrow *bool*

interpret *c:complete-space dist uniformity ?open*

proof –

have $\exists d.$ *Metric-space (UNIV :: ereal set) d* \wedge

```

      Metric-space.mtopology UNIV d = euclidean ∧
      Metric-space.mcomplete UNIV d
    by (metis Polish-space-ereal Metric-space.topspace-mtopology Polish-space-def
        completely-metrizable-space-def topspace-euclidean)
    hence Metric-space (UNIV :: ereal set) dist ∧
      Metric-space.mtopology (UNIV :: ereal set) dist = euclidean ∧
      Metric-space.mcomplete (UNIV :: ereal set) dist
    unfolding dist-ereal-def by(rule someI-ex)
    with completely-metrizable-space-metric-space show class.complete-space dist
    uniformity ?open
      by(fastforce simp: uniformity-ereal-def)
    qed
    have [simp]:topological-space.convergent ?open = convergent
    proof
      fix x :: nat ⇒ ereal
      have *:class.topological-space ?open
        by standard auto
      show topological-space.convergent open x = convergent x
        by(simp add: topological-space.convergent-def topological-space.nhds-def * con-
            vergent-def nhds-def)
    qed
    have [simp]:uniform-space.Cauchy (uniformity :: (ereal × ereal) filter) = Cauchy
      by(auto simp add: metric-space.Cauchy-def[OF metric-space-axioms] Cauchy-def)
    fix x :: nat ⇒ ereal
    show Cauchy x ⇒ convergent x
      using c.Cauchy-convergent by(auto simp: Cauchy-def)
    qed
  end
end

```

instantiation *ennreal* :: *metric-space*
begin

definition *dist-ennreal* :: *ennreal* ⇒ *ennreal* ⇒ *real*
 where *dist-ennreal* ≡ SOME d. Metric-space UNIV d ∧
 Metric-space.mtopology UNIV d = euclidean ∧
 Metric-space.mcomplete UNIV d

definition *uniformity-ennreal* :: (*ennreal* × *ennreal*) filter
 where *uniformity-ennreal* ≡ $\prod_{e \in \{0 < ..\}}$. principal $\{(x, y). \text{dist } x \ y < e\}$

instance

```

proof –
  let ?open = open :: ennreal set ⇒ bool
  interpret c:complete-space dist uniformity ?open
  proof –
    have ∃ d. Metric-space (UNIV :: ennreal set) d ∧
      Metric-space.mtopology UNIV d = euclidean ∧
      Metric-space.mcomplete UNIV d

```

```

    by (metis Polish-space-ennreal Metric-space.topspace-mtopology Polish-space-def
        completely-metrizable-space-def topspace-euclidean)
    hence Metric-space (UNIV :: ennreal set) dist ∧
        Metric-space.mtopology (UNIV :: ennreal set) dist = euclidean ∧
        Metric-space.mcomplete (UNIV :: ennreal set) dist
    unfolding dist-ennreal-def by(rule someI-ex)
    with completely-metrizable-space-metric-space show class.complete-space dist
uniformity ?open
    by(fastforce simp: uniformity-ennreal-def)
qed
have [simp]:topological-space.convergent ?open = convergent
proof
  fix x :: nat ⇒ ennreal
  have *:class.topological-space ?open
    by standard auto
  show topological-space.convergent open x = convergent x
  by(simp add: topological-space.convergent-def[OF *] topological-space.nhds-def[OF
*] convergent-def nhds-def)
qed
show OFCLASS(ennreal, metric-space-class)
  by standard (use uniformity-ennreal-def c.open-uniformity c.dist-triangle2 c.Cauchy-convergent
in auto)
qed

end

instantiation ennreal :: polish-space
begin

instance
proof
  let ?open = open :: ennreal set ⇒ bool
  interpret c:complete-space dist uniformity ?open
  proof -
    have ∃ d. Metric-space (UNIV :: ennreal set) d ∧
        Metric-space.mtopology UNIV d = euclidean ∧
        Metric-space.mcomplete UNIV d
    by (metis Polish-space-ennreal Metric-space.topspace-mtopology Polish-space-def
        completely-metrizable-space-def topspace-euclidean)
    hence Metric-space (UNIV :: ennreal set) dist ∧
        Metric-space.mtopology (UNIV :: ennreal set) dist = euclidean ∧
        Metric-space.mcomplete (UNIV :: ennreal set) dist
    unfolding dist-ennreal-def by(rule someI-ex)
    with completely-metrizable-space-metric-space show class.complete-space dist
uniformity ?open
    by(fastforce simp: uniformity-ennreal-def)
  qed
  have [simp]:topological-space.convergent ?open = convergent
  proof

```

```

fix x :: nat ⇒ ennreal
have *:class.topological-space ?open
  by standard auto
show topological-space.convergent open x = convergent x
  by(simp add: topological-space.convergent-def topological-space.nhds-def * con-
vergent-def nhds-def)
qed
have [simp]:uniform-space.Cauchy (uniformity :: (ennreal × ennreal) filter) =
Cauchy
  by(auto simp add: metric-space.Cauchy-def[OF metric-space-axioms] Cauchy-def)
fix x :: nat ⇒ ennreal
show Cauchy x ⇒ convergent x
  using c.Cauchy-convergent by(auto simp: Cauchy-def)
qed

end

```

1.2 Lemmas for Infinite Sum

```

lemma uniformly-continuous-add-ennreal: isUCont (λ(x::ennreal, y). x + y)
proof(safe intro!: compact-uniformly-continuous)
  have compact (UNIV × UNIV :: (ennreal × ennreal) set)
    by(intro compact-Times compact-UNIV)
  thus compact (UNIV :: (ennreal × ennreal) set)
    by simp
qed(auto intro!: continuous-on-add-ennreal continuous-on-fst continuous-on-snd simp:
split-beta')

```

```

lemma infsum-eq-suminf:
  assumes f summable-on UNIV
  shows (∑∞ n∈UNIV. f n) = suminf f
  using has-sum-imp-sums[OF has-sum-infsum[OF assms]]
  by (simp add: sums-iff)

```

```

lemma infsum-Sigma-ennreal:
  fixes f :: - ⇒ ennreal
  shows infsum f (Sigma A B) = infsum (λx. infsum (λy. f (x, y)) (B x)) A
  by(auto intro!: uniformly-continuous-add-ennreal infsum-Sigma nonneg-summable-on-complete)

```

```

lemma infsum-swap-ennreal:
  fixes f :: - ⇒ - ⇒ ennreal
  shows infsum (λx. infsum (λy. f x y) B) A = infsum (λy. infsum (λx. f x y) A)
B
  by(auto intro!: infsum-swap uniformly-continuous-add-ennreal nonneg-summable-on-complete)

```

```

lemma has-sum-cmult-right-ennreal:
  fixes f :: - ⇒ ennreal
  assumes c < ⊤ (f has-sum a) A
  shows ((λx. c * f x) has-sum c * a) A

```

using *ennreal-tendsto-cmult*[*OF assms(1)*] *assms(2)*
 by (*force simp add: has-sum-def sum-distrib-left*)

lemma *infsum-cmult-right-ennreal*:

fixes $f :: - \Rightarrow \text{ennreal}$

assumes $c < \top$

shows $(\sum_{\infty} x \in A. c * f x) = c * \text{infsum } f A$

by (*simp add: assms has-sum-cmult-right-ennreal infsumI nonneg-summable-on-complete*)

lemma *ennreal-sum-SUP-eq*:

fixes $f :: \text{nat} \Rightarrow - \Rightarrow \text{ennreal}$

assumes $\text{finite } A \wedge x. x \in A \implies \text{incseq } (\lambda j. f j x)$

shows $(\sum_{i \in A}. \bigsqcup n. f n i) = (\bigsqcup n. \sum_{i \in A}. f n i)$

using *assms*

proof *induction*

case *empty*

then show *?case*

by *simp*

next

case *ih:(insert x F)*

show *?case (is ?lhs = ?rhs)*

proof –

have $?lhs = (\bigsqcup n. f n x) + (\bigsqcup n. \text{sum } (f n) F)$

using *ih by simp*

also have $\dots = (\bigsqcup n. f n x + \text{sum } (f n) F)$

using *ih by (auto intro!: incseq-sumI2 ennreal-SUP-add[symmetric])*

also have $\dots = ?rhs$

using *ih by simp*

finally show *?thesis* .

qed

qed

lemma *ennreal-infsum-Sup-eq*:

fixes $f :: \text{nat} \Rightarrow - \Rightarrow \text{ennreal}$

assumes $\wedge x. x \in A \implies \text{incseq } (\lambda j. f j x)$

shows $(\sum_{\infty} x \in A. (\text{SUP } j. f j x)) = (\text{SUP } j. (\sum_{\infty} x \in A. f j x))$ (**is** *?lhs = ?rhs*)

proof –

have $?lhs = (\bigsqcup (\text{sum } (\lambda x. \bigsqcup j. f j x) \text{ ‘ } \{F. \text{finite } F \wedge F \subseteq A\}))$

by(*auto intro!: nonneg-infsum-complete simp: SUP-upper2 assms*)

also have $\dots = (\bigsqcup A \in \{F. \text{finite } F \wedge F \subseteq A\}. \bigsqcup j. \text{sum } (f j) A)$

using *assms by (auto intro!: SUP-cong ennreal-sum-SUP-eq)*

also have $\dots = (\bigsqcup j. \bigsqcup A \in \{F. \text{finite } F \wedge F \subseteq A\}. \text{sum } (f j) A)$

using *SUP-commute by fast*

also have $\dots = ?rhs$

by(*subst nonneg-infsum-complete*) (*use assms in auto*)

finally show *?thesis* .

qed

lemma *bounded-infsum-summable*:

assumes $\bigwedge x. x \in A \implies f x \geq 0$ $(\sum_{\infty x \in A. \text{ennreal } (f x)} < \text{top})$
shows f summable-on A
proof(rule nonneg-bdd-above-summable-on)
from $\text{assms}(2)$ **obtain** K **where** $K: (\sum_{\infty x \in A. \text{ennreal } (f x)} \leq \text{ennreal } K \ K \geq 0)$
using less-top-ennreal **by** force
show $\text{bdd-above } (\text{sum } f \text{ ' } \{F. F \subseteq A \wedge \text{finite } F\})$
proof(safe intro!: $\text{bdd-aboveI}[\text{where } M=K]$)
fix A'
assume $A': A' \subseteq A$ $\text{finite } A'$
have $(\sum_{\infty x \in A. \text{ennreal } (f x)} = (\bigsqcup (\text{sum } (\lambda x. \text{ennreal } (f x)) \text{ ' } \{F. \text{finite } F \wedge F \subseteq A\}))$
by ($\text{simp add: nonneg-infsum-complete}$)
also have $\dots = (\bigsqcup ((\lambda F. \text{ennreal } (\text{sum } f F)) \text{ ' } \{F. \text{finite } F \wedge F \subseteq A\}))$
by($\text{auto intro!: SUP-cong sum-ennreal assms}$)
finally have $(\bigsqcup ((\lambda F. \text{ennreal } (\text{sum } f F)) \text{ ' } \{F. \text{finite } F \wedge F \subseteq A\})) \leq \text{ennreal } K$
using K **by** order
hence $\text{ennreal } (\text{sum } f A') \leq \text{ennreal } K$
by ($\text{simp add: } A' \text{ SUP-le-iff}$)
thus $\text{sum } f A' \leq K$
by ($\text{simp add: } K(2)$)
qed
qed fact

lemma $\text{infsum-less-top-dest}$:

fixes $f :: - \Rightarrow - :: \{\text{ordered-comm-monoid-add, topological-comm-monoid-add, } t2\text{-space, complete-linorder, linorder-topology}\}$
assumes $(\sum_{\infty x \in A. f x} < \text{top}) \wedge x. x \in A \implies f x \geq 0 \ x \in A$
shows $f x < \text{top}$
proof(rule ccontr)
assume $f: \neg f x < \text{top}$
have $(\sum_{\infty x \in A. f x} = (\sum_{\infty y \in A - \{x\} \cup \{x\}. f y})$
by(rule $\text{arg-cong}[\text{where } f=\text{infsum } -]$) (use assms in auto)
also have $\dots = (\sum_{\infty y \in A - \{x\}. f y} + (\sum_{\infty y \in \{x\}. f y})$
using $\text{assms}(2)$ **by**($\text{intro infsum-Un-disjoint}$) ($\text{auto intro!: nonneg-summable-on-complete}$)
also have $\dots = (\sum_{\infty y \in A - \{x\}. f y} + \text{top})$
using $f \text{ top.not-eq-extremum}$ **by** fastforce
also have $\dots = \text{top}$
by($\text{auto intro!: add-top infsum-nonneg assms}$)
finally show False
using $\text{assms}(1)$ **by** simp
qed

lemma infsum-ennreal-eq :

assumes f summable-on A $\bigwedge x. x \in A \implies f x \geq 0$
shows $(\sum_{\infty x \in A. \text{ennreal } (f x)} = \text{ennreal } (\sum_{\infty x \in A. f x})$
proof –
have $(\sum_{\infty x \in A. \text{ennreal } (f x)} = (\bigsqcup (\text{sum } (\lambda x. \text{ennreal } (f x)) \text{ ' } \{F. \text{finite } F \wedge$

$F \subseteq A$))
 by (simp add: nonneg-infsum-complete)
 also have ... = ($\bigsqcup ((\lambda F. \text{ennreal } (\text{sum } f F)) \text{ ` } \{F. \text{finite } F \wedge F \subseteq A\})$)
 by(auto intro!: SUP-cong sum-ennreal assms)
 also have ... = $\text{ennreal } (\sum_{\infty} x \in A. f x)$
 using infsum-nonneg-is-SUPREMUM-ennreal[OF assms] by simp
 finally show ?thesis .
 qed

lemma *abs-summable-on-integrable-iff*:
 fixes $f :: - \Rightarrow - :: \{\text{banach, second-countable-topology}\}$
 shows $\text{Infinite-Sum.abs-summable-on } f A \longleftrightarrow \text{integrable } (\text{count-space } A) f$
 by (simp add: abs-summable-equivalent abs-summable-on-def)

lemma *infsum-eq-integral*:
 fixes $f :: - \Rightarrow - :: \{\text{banach, second-countable-topology}\}$
 assumes $\text{Infinite-Sum.abs-summable-on } f A$
 shows $\text{infsum } f A = \text{integral}^L (\text{count-space } A) f$
 using assms infsetsum-infsum[of $f A$, symmetric]
 by(auto simp: abs-summable-on-integrable-iff abs-summable-on-def infsetsum-def)

end

theory *Coproduct-Measure*
 imports *Lemmas-Coproduct-Measure*
 HOL-Analysis.Analysis
 begin

2 Binary Coproduct Measures

definition *copair-measure* :: [$'a$ measure, $'b$ measure] \Rightarrow ($'a + 'b$) measure (**infixr** \oplus_M 65) **where**
 $M \oplus_M N = \text{measure-of } (\text{space } M \langle + \rangle \text{ space } N)$
 ($\{\text{Inl } 'A \mid A. A \in \text{sets } M\} \cup \{\text{Inr } 'A \mid A. A \in \text{sets } N\}$)
 ($\lambda A. \text{emeasure } M (\text{Inl } - 'A) + \text{emeasure } N (\text{Inr } - 'A)$)

2.1 The Measurable Space and Measurability

lemma
 shows *space-copair-measure*: $\text{space } (\text{copair-measure } M N) = \text{space } M \langle + \rangle \text{ space } N$
and *sets-copair-measure-sigma*:
 $\text{sets } (\text{copair-measure } M N)$
 = $\text{sigma-sets } (\text{space } M \langle + \rangle \text{ space } N) (\{\text{Inl } 'A \mid A. A \in \text{sets } M\} \cup \{\text{Inr } 'A \mid A. A \in \text{sets } N\})$
and *Inl-measurable*[measurable]: $\text{Inl} \in M \rightarrow_M M \oplus_M N$
and *Inr-measurable*[measurable]: $\text{Inr} \in N \rightarrow_M M \oplus_M N$
proof –

have $1: (\{Inl \text{ ' } A \mid A. A \in \text{sets } M\} \cup \{Inr \text{ ' } A \mid A. A \in \text{sets } N\}) \subseteq Pow (\text{space } M \lt + \gt \text{space } N)$
using $\text{sets.sets-into-space}[of - M] \text{sets.sets-into-space}[of - N]$ **by** fastforce
show $\text{space } (\text{copair-measure } M N) = \text{space } M \lt + \gt \text{space } N$
and $2: \text{sets } (\text{copair-measure } M N)$
 $= \text{sigma-sets } (\text{space } M \lt + \gt \text{space } N) (\{Inl \text{ ' } A \mid A. A \in \text{sets } M\} \cup \{Inr \text{ ' } A \mid A. A \in \text{sets } N\})$
by $(\text{simp-all add: copair-measure-def sets-measure-of}[OF 1] \text{space-measure-of}[OF 1])$
show $Inl \in M \rightarrow_M M \oplus_M N \text{Inr} \in N \rightarrow_M M \oplus_M N$
by $(\text{auto intro!: measurable-sigma-sets}[OF 2 1] \text{simp: vimage-def image-def})$
qed

lemma $\text{sets-copair-measure-cong}$:
 $\text{sets } M1 = \text{sets } M2 \implies \text{sets } N1 = \text{sets } N2 \implies \text{sets } (M1 \oplus_M N1) = \text{sets } (M2 \oplus_M N2)$
by $(\text{simp cong: sets-eq-imp-space-eq add: sets-copair-measure-sigma})$

lemma $\text{measurable-image-Inl}[\text{measurable}]$: $A \in \text{sets } M \implies Inl \text{ ' } A \in \text{sets } (M \oplus_M N)$
using $\text{sets-copair-measure-sigma}$ **by** fastforce

lemma $\text{measurable-image-Inr}[\text{measurable}]$: $A \in \text{sets } N \implies Inr \text{ ' } A \in \text{sets } (M \oplus_M N)$
using $\text{sets-copair-measure-sigma}$ **by** fastforce

lemma $\text{measurable-vimage-Inl}$:
assumes $[\text{measurable}]$: $A \in \text{sets } (M \oplus_M N)$
shows $Inl \text{ - ' } A \in \text{sets } M$
proof –
have $Inl \text{ - ' } A = Inl \text{ - ' } A \cap \text{space } M$
using $\text{sets.sets-into-space}[OF \text{ assms}]$
by $(\text{auto simp add: space-copair-measure})$
also have $\dots \in \text{sets } M$
by simp
finally show $?thesis .$
qed

lemma $\text{measurable-vimage-Inr}$:
assumes $[\text{measurable}]$: $A \in \text{sets } (M \oplus_M N)$
shows $Inr \text{ - ' } A \in \text{sets } N$
proof –
have $Inr \text{ - ' } A = Inr \text{ - ' } A \cap \text{space } N$
using $\text{sets.sets-into-space}[OF \text{ assms}]$
by $(\text{auto simp add: space-copair-measure})$
also have $\dots \in \text{sets } N$
by simp
finally show $?thesis .$
qed

lemma *in-sets-copair-measure-iff*:

$A \in \text{sets } (\text{copair-measure } M \ N) \iff \text{Inl } -' A \in \text{sets } M \wedge \text{Inr } -' A \in \text{sets } N$

proof *safe*

assume [*measurable*]: $\text{Inl } -' A \in \text{sets } M \ \text{Inr } -' A \in \text{sets } N$

have $A = ((\text{Inl } -' \text{Inl } -' A) \cup (\text{Inr } -' \text{Inr } -' A))$

by (*simp add: vimage-def image-def*) (*safe, metis obj-sumE*)

also have $\dots \in \text{sets } (\text{copair-measure } M \ N)$

by *measurable*

finally show $A \in \text{sets } (\text{copair-measure } M \ N)$.

qed(*use measurable-vimage-Inl measurable-vimage-Inr in auto*)

lemma *measurable-copair-Inl-Inr*:

assumes [*measurable*]: $(\lambda x. f (\text{Inl } x)) \in M \rightarrow_M L \ (\lambda x. f (\text{Inr } x)) \in N \rightarrow_M L$

shows $f \in M \oplus_M N \rightarrow_M L$

proof(*rule measurableI*)

fix A

assume [*measurable*]: $A \in \text{sets } L$

have $f -' A = \text{Inl } -' ((\lambda x. f (\text{Inl } x)) -' A) \cup \text{Inr } -' ((\lambda x. f (\text{Inr } x)) -' A)$

by (*simp add: image-def vimage-def*) (*safe, metis obj-sumE*)

hence $f -' A \cap \text{space } (M \oplus_M N)$

$= \text{Inl } -' ((\lambda x. f (\text{Inl } x)) -' A \cap \text{space } M) \cup \text{Inr } -' ((\lambda x. f (\text{Inr } x)) -' A \cap$

$\text{space } N)$

by(*auto simp: space-copair-measure*)

also have $\dots \in \text{sets } (M \oplus_M N)$

by *measurable*

finally show $f -' A \cap \text{space } (M \oplus_M N) \in \text{sets } (M \oplus_M N)$.

next

show $\bigwedge x. x \in \text{space } (M \oplus_M N) \implies f x \in \text{space } L$

using *measurable-space[OF assms(1)] measurable-space[OF assms(2)]*

by(*auto simp add: space-copair-measure*)

qed

corollary *measurable-copair-measure-iff*:

$f \in M \oplus_M N \rightarrow_M L \iff (\lambda x. f (\text{Inl } x)) \in M \rightarrow_M L \wedge (\lambda x. f (\text{Inr } x)) \in N \rightarrow_M L$

by(*auto simp add: measurable-copair-Inl-Inr*)

lemma *measurable-copair-dest1*:

assumes [*measurable*]: $f \in L \rightarrow_M M \oplus_M N$ **and** $f -' (\text{Inl } -' \text{space } M) \cap \text{space } L = \text{space } L$

obtains f' **where** $f' \in L \rightarrow_M M \ \bigwedge x. x \in \text{space } L \implies f x = \text{Inl } (f' x)$

proof -

define f' **where** $f' \equiv (\lambda x. \text{SOME } y. f x = \text{Inl } y)$

have $f': \bigwedge x. x \in \text{space } L \implies f x = \text{Inl } (f' x)$

unfolding f' -*def* **by**(*rule someI-ex*) (*use assms(2) in blast*)

moreover have $f' \in L \rightarrow_M M$

proof(*rule measurableI*)

show $\bigwedge x. x \in \text{space } L \implies f' x \in \text{space } M$

```

    using f' measurable-space[OF assms(1)]
    by(auto simp: space-copair-measure)
next
fix A
assume A[measurable]:A ∈ sets M
have [simp]:f' -' A ∩ space L = f -' (Inl ' A) ∩ space L
    using f' sets.sets-into-space[OF A] by auto
show f' -' A ∩ space L ∈ sets L
    by auto
qed
ultimately show ?thesis
    using that by blast
qed

lemma measurable-copair-dest2:
  assumes [measurable]:f ∈ L →M M ⊕M N and f -' (Inr ' space N) ∩ space
L = space L
  obtains f' where f' ∈ L →M N ∧x. x ∈ space L ⇒ f x = Inr (f' x)
proof -
  define f' where f' ≡ (λx. SOME y. f x = Inr y)
  have f':∧x. x ∈ space L ⇒ f x = Inr (f' x)
    unfolding f'-def by(rule someI-ex) (use assms(2) in blast)
  moreover have f' ∈ L →M N
  proof(rule measurableI)
    show ∧x. x ∈ space L ⇒ f' x ∈ space N
      using f' measurable-space[OF assms(1)]
      by(auto simp: space-copair-measure)
  next
  fix A
  assume A[measurable]:A ∈ sets N
  have [simp]:f' -' A ∩ space L = f -' (Inr ' A) ∩ space L
      using f' sets.sets-into-space[OF A] by auto
  show f' -' A ∩ space L ∈ sets L
      by auto
  qed
  ultimately show ?thesis
      using that by blast
  qed

lemma measurable-copair-dest3:
  assumes [measurable]:f ∈ L →M M ⊕M N
  and f -' (Inl ' space M) ∩ space L ⊂ space L f -' (Inr ' space N) ∩ space L
  ⊂ space L
  obtains f' f'' where f' ∈ L →M M f'' ∈ L →M N
  ∧x. x ∈ space L ⇒ x ∈ f -' Inl ' space M ⇒ f x = Inl (f' x)
  ∧x. x ∈ space L ⇒ x ∉ f -' Inl ' space M ⇒ f x = Inr (f'' x)
proof -
  have ne:space M ≠ {} space N ≠ {}
    using assms(2,3) measurable-space[OF assms(1)] by(fastforce simp: space-copair-measure)+

```

```

define  $m$  where  $m \equiv \text{SOME } y. y \in \text{space } M$ 
define  $n$  where  $n \equiv \text{SOME } y. y \in \text{space } N$ 
have  $m[\text{measurable}, \text{simp}]: m \in \text{space } M$  and  $n[\text{measurable}, \text{simp}]: n \in \text{space } N$ 
using  $ne$  by( $\text{auto simp: } n\text{-def } m\text{-def some-in-eq}$ )
define  $f'$  where  $f' \equiv (\lambda x. \text{if } x \in f - ' \text{Inl } ' \text{space } M \text{ then } \text{SOME } y. f x = \text{Inl } y$ 
else } m)
have  $\bigwedge x. x \in \text{space } L \implies x \in f - ' \text{Inl } ' \text{space } M \implies f x = \text{Inl } (\text{SOME } y. f x$ 
 $= \text{Inl } y)$ 
unfolding  $f'\text{-def}$  by( $\text{rule someI-ex}$ ) ( $\text{use } \text{assms}(2)$  in  $\text{blast}$ )
hence  $f': \bigwedge x. x \in \text{space } L \implies x \in f - ' \text{Inl } ' \text{space } M \implies f x = \text{Inl } (f' x)$ 
by( $\text{simp add: } f'\text{-def}$ )
hence  $f'\text{-space}: x \in \text{space } L \implies f' x \in \text{space } M$  for  $x$ 
using  $\text{measurable-space}[OF \text{ assms}(1)]$ 
by( $\text{cases } x \in f - ' \text{Inl } ' \text{space } M$ ) ( $\text{auto simp: space-copair-measure } f'\text{-def}$ )
define  $f''$  where  $f'' \equiv (\lambda x. \text{if } x \notin f - ' \text{Inl } ' \text{space } M \text{ then } \text{SOME } y. f x = \text{Inr } y$ 
else } n)
have  $*$ :  $\bigwedge x. x \in \text{space } L \implies x \notin f - ' \text{Inl } ' \text{space } M \implies x \in f - ' \text{Inr } ' \text{space } N$ 
using  $\text{measurable-space}[OF \text{ assms}(1)]$  by( $\text{fastforce simp: space-copair-measure}$ )
have  $\bigwedge x. x \in \text{space } L \implies x \notin f - ' \text{Inl } ' \text{space } M \implies f x = \text{Inr } (\text{SOME } y. f x$ 
 $= \text{Inr } y)$ 
unfolding  $f''\text{-def}$  by( $\text{rule someI-ex}$ ) ( $\text{use } *$  in  $\text{blast}$ )
hence  $f'': \bigwedge x. x \in \text{space } L \implies x \notin f - ' \text{Inl } ' \text{space } M \implies f x = \text{Inr } (f'' x)$ 
by( $\text{simp add: } f''\text{-def}$ )
hence  $f''\text{-space}: x \in \text{space } L \implies f'' x \in \text{space } N$  for  $x$ 
using  $\text{measurable-space}[OF \text{ assms}(1), \text{of } x]$ 
by( $\text{cases } x \notin f - ' \text{Inl } ' \text{space } M$ ) ( $\text{auto simp add: space-copair-measure } f''\text{-def}$ )
have  $f' \in L \rightarrow_M M$ 
proof -
have  $f' = (\lambda x. \text{if } x \in f - ' \text{Inl } ' \text{space } M \text{ then } f' x \text{ else } m)$ 
by( $\text{auto simp add: } f'\text{-def}$ )
also have  $\dots \in L \rightarrow_M M$ 
proof( $\text{intro measurable-restrict-space-iff}[THEN \text{ iffD1}]$   $\text{measurableI}$ )
fix  $A$ 
assume  $A[\text{measurable}]: A \in \text{sets } M$ 
have  $[\text{measurable}]: f \in \text{restrict-space } L (f - ' \text{Inl } ' \text{space } M) \rightarrow_M M \oplus_M N$ 
by( $\text{auto intro!: measurable-restrict-space1}$ )
have  $[\text{simp}]: f' - ' A \cap \text{space } (\text{restrict-space } L (f - ' \text{Inl } ' \text{space } M))$ 
 $= f - ' (\text{Inl } ' A) \cap \text{space } (\text{restrict-space } L (f - ' \text{Inl } ' \text{space } M))$ 
using  $f' \text{ sets.sets-into-space}[OF A]$  by( $\text{fastforce simp: space-restrict-space}$ )
show  $f' - ' A \cap \text{space } (\text{restrict-space } L (f - ' \text{Inl } ' \text{space } M))$ 
 $\in \text{sets } (\text{restrict-space } L (f - ' \text{Inl } ' \text{space } M))$ 
by  $\text{simp}$ 
next
show  $\bigwedge x. x \in \text{space } (\text{restrict-space } L (f - ' \text{Inl } ' \text{space } M)) \implies f' x \in \text{space } M$ 
by( $\text{auto simp: space-restrict-space } f'\text{-space}$ )
qed  $\text{simp-all}$ 
finally show  $?thesis$  .
qed

```

```

moreover have  $f'' \in L \rightarrow_M N$ 
proof -
  have  $f'' = (\lambda x. \text{if } x \notin f - ' \text{Inl } ' \text{ space } M \text{ then } f'' x \text{ else } n)$ 
  by(auto simp add: f''-def)
  also have  $\dots \in L \rightarrow_M N$ 
  proof(rule measurable-If-restrict-space-iff[THEN iffD2, OF - conjI[OF measurableI]])
    fix  $A$ 
    assume  $A[\text{measurable}]: A \in \text{sets } N$ 
    have  $f: f \in \text{restrict-space } L \{x. x \notin f - ' \text{Inl } ' \text{ space } M\} \rightarrow_M M \oplus_M N$ 
    by(auto intro!: measurable-restrict-space1)
    have  $1: f'' - ' A \cap \text{space } (\text{restrict-space } L \{x. x \notin f - ' \text{Inl } ' \text{ space } M\})$ 
       $= f - ' (\text{Inr } ' A) \cap \text{space } (\text{restrict-space } L \{x. x \notin f - ' \text{Inl } ' \text{ space } M\})$ 
    using  $f'' \text{ sets.sets-into-space}[OF A]$  by(fastforce simp: space-restrict-space)
    show  $f'' - ' A \cap \text{space } (\text{restrict-space } L \{x. x \notin f - ' \text{Inl } ' \text{ space } M\})$ 
       $\in \text{sets } (\text{restrict-space } L \{x. x \notin f - ' \text{Inl } ' \text{ space } M\})$ 
    unfolding 1 using f by simp
  next
    show  $\bigwedge x. x \in \text{space } (\text{restrict-space } L \{x. x \notin f - ' \text{Inl } ' \text{ space } M\}) \implies f'' x \in \text{space } N$ 
    by(auto simp: space-restrict-space f''-space)
  qed simp-all
  finally show ?thesis .
qed
ultimately show ?thesis
using that f' f'' by blast
qed

```

2.2 Measures

lemma *emeasure-copair-measure*:

```

assumes [measurable]:  $A \in \text{sets } (M \oplus_M N)$ 
shows  $\text{emeasure } (M \oplus_M N) A = \text{emeasure } M (\text{Inl } - ' A) + \text{emeasure } N (\text{Inr } - ' A)$ 
proof(rule emeasure-measure-of)
  show  $\{\text{Inl } ' A \mid A. A \in \text{sets } M\} \cup \{\text{Inr } ' A \mid A. A \in \text{sets } N\} \subseteq \text{Pow } (\text{space } M \langle + \rangle \text{ space } N)$ 
  using sets.sets-into-space[of - M] sets.sets-into-space[of - N] by fastforce
  show  $A \in \text{sets } (M \oplus_M N)$ 
  by fact
  show countably-additive (sets  $(M \oplus_M N)$ ) ( $\lambda a. \text{emeasure } M (\text{Inl } - ' a) + \text{emeasure } N (\text{Inr } - ' a)$ )
  proof(safe intro!: countably-additiveI)
    note [measurable] = measurable-vimage-Inl[of - M N] measurable-vimage-Inr[of - M N]
    fix  $A :: \text{nat} \Rightarrow - \text{ set}$ 
    assume  $h: \text{range } A \subseteq \text{sets } (M \oplus_M N)$  disjoint-family A
    then have [measurable]:  $\bigwedge i. A i \in \text{sets } (M \oplus_M N)$ 
    by blast

```

```

have disj:disjoint-family ( $\lambda i. \text{Inl } -' A i$ ) disjoint-family ( $\lambda i. \text{Inr } -' A i$ )
  using h by(auto simp: disjoint-family-on-def)
show ( $\sum i. \text{emeasure } M (\text{Inl } -' A i) + \text{emeasure } N (\text{Inr } -' A i)$ )
  =  $\text{emeasure } M (\text{Inl } -' \bigcup (\text{range } A)) + \text{emeasure } N (\text{Inr } -' \bigcup (\text{range } A))$ 
is ?lhs = ?rhs)
proof -
  have ?lhs = ( $\sum i. \text{emeasure } M (\text{Inl } -' A i) + \sum i. \text{emeasure } N (\text{Inr } -' A i)$ )
  by(simp add: suminf-add)
  also have ... =  $\text{emeasure } M (\bigcup i. (\text{Inl } -' A i)) + \text{emeasure } N (\bigcup i. (\text{Inr } -' A i))$ 
  proof -
    have ( $\sum i. \text{emeasure } M (\text{Inl } -' A i) = \text{emeasure } M (\bigcup i. (\text{Inl } -' A i))$ )
      ( $\sum i. \text{emeasure } N (\text{Inr } -' A i) = \text{emeasure } N (\bigcup i. (\text{Inr } -' A i))$ )
      by(auto intro!: suminf-emeasure disj)
    thus ?thesis
    by argo
  qed
  also have ... = ?rhs
  by(simp add: vimage-UN)
  finally show ?thesis .
qed
qed
qed(auto simp: positive-def copair-measure-def)

```

```

lemma emeasure-copair-measure-space:
   $\text{emeasure } (M \oplus_M N) (\text{space } (M \oplus_M N)) = \text{emeasure } M (\text{space } M) + \text{emeasure } N (\text{space } N)$ 
proof -
  have [simp]: $\text{Inl } -' \text{space } (M \oplus_M N) = \text{space } M$   $\text{Inr } -' \text{space } (M \oplus_M N) = \text{space } N$ 
  by(auto simp: space-copair-measure)
  show ?thesis
  by(simp add: emeasure-copair-measure)
qed

```

```

corollary
  shows emeasure-copair-measure-Inl:  $A \in \text{sets } M \implies \text{emeasure } (M \oplus_M N) (\text{Inl } -' A) = \text{emeasure } M A$ 
  and emeasure-copair-measure-Inr:  $B \in \text{sets } N \implies \text{emeasure } (M \oplus_M N) (\text{Inr } -' B) = \text{emeasure } N B$ 
proof -
  have [simp]: $\text{Inl } -' \text{Inl } -' A = A$   $\text{Inr } -' \text{Inl } -' A = \{\}$   $\text{Inl } -' \text{Inr } -' B = \{\}$   $\text{Inr } -' \text{Inr } -' B = B$ 
  by auto
  show  $A \in \text{sets } M \implies \text{emeasure } (M \oplus_M N) (\text{Inl } -' A) = \text{emeasure } M A$ 
     $B \in \text{sets } N \implies \text{emeasure } (M \oplus_M N) (\text{Inr } -' B) = \text{emeasure } N B$ 
  by(simp-all add: emeasure-copair-measure)
qed

```


lemma *measure-copair-measure*:

assumes $[measurable]: A \in \text{sets } (M \oplus_M N)$ *emeasure* $(M \oplus_M N) A < \infty$
shows $\text{measure } (M \oplus_M N) A = \text{measure } M (\text{Inl } - ' A) + \text{measure } N (\text{Inr } - ' A)$
using *assms(2)* **by** $(\text{auto simp add: } \text{emeasure-copair-measure measure-def intro!: } \text{enn2real-plus})$

lemma

shows *measure-copair-measure-Inl*: $A \in \text{sets } M \implies \text{measure } (M \oplus_M N) (\text{Inl } ' A) = \text{measure } M A$
and *measure-copair-measure-Inr*: $B \in \text{sets } N \implies \text{measure } (M \oplus_M N) (\text{Inr } ' B) = \text{measure } N B$
by $(\text{auto simp: } \text{emeasure-copair-measure-Inl measure-def } \text{emeasure-copair-measure-Inr})$

2.3 Finiteness

lemma *finite-measure-copair-measure*: $\text{finite-measure } M \implies \text{finite-measure } N \implies \text{finite-measure } (M \oplus_M N)$

by $(\text{auto intro!: } \text{finite-measureI simp: } \text{emeasure-copair-measure-space finite-measure.finite-emeasure-space})$

2.4 σ -Finiteness

lemma *sigma-finite-measure-copair-measure*:

assumes *sigma-finite-measure* M *sigma-finite-measure* N

shows *sigma-finite-measure* $(M \oplus_M N)$

proof –

obtain $A B$ **where** $AB[measurable]: \bigwedge i. A i \in \text{sets } M (\bigcup (\text{range } A)) = \text{space } M \bigwedge i::\text{nat. } \text{emeasure } M (A i) \neq \infty$

$\bigwedge i. B i \in \text{sets } N (\bigcup (\text{range } B)) = \text{space } N \bigwedge i::\text{nat. } \text{emeasure } N (B i) \neq \infty$

by $(\text{metis range-subsetD } \text{sigma-finite-measure.sigma-finite assms})$

then have $*(\bigcup (\text{range } (\lambda i. \text{Inl } ' (A i) \cup \text{Inr } ' (B i)))) = \text{space } (M \oplus_M N)$

unfolding *space-copair-measure Plus-def* **by** *fastforce*

have $[simp]: \bigwedge i. \text{Inl } - ' \text{Inl } ' A i \cup \text{Inl } - ' \text{Inr } ' B i = A i \bigwedge i. \text{Inr } - ' \text{Inl } ' A i \cup \text{Inr } - ' \text{Inr } ' B i = B i$

using *sets.sets-into-space AB(1,4)* **by** *blast+*

show *?thesis*

apply *standard*

using $AB *$ **by** $(\text{auto intro!: } \text{exI}[\text{where } x = \text{range } (\lambda i. \text{Inl } ' (A i) \cup \text{Inr } ' (B i))])$
simp: space-copair-measure emeasure-copair-measure)

qed

2.5 Non-Negative Integral

lemma *nn-integral-copair-measure*:

assumes $f \in \text{borel-measurable } (M \oplus_M N)$

shows $(\int^{+x}. f x \partial(M \oplus_M N)) = (\int^{+x}. f (\text{Inl } x) \partial M) + (\int^{+x}. f (\text{Inr } x) \partial N)$

using *assms*

proof *induction*

case $(\text{cong } f g)$

moreover hence $\bigwedge x. x \in \text{space } M \implies f (\text{Inl } x) = g (\text{Inl } x)$
 $\bigwedge x. x \in \text{space } N \implies f (\text{Inr } x) = g (\text{Inr } x)$
by(*auto simp: space-copair-measure*)
ultimately show *?case*
by(*simp cong: nn-integral-cong*)
next
case [*measurable*]:(*set A*)
note [*measurable*] = *measurable-vimage-Inl[of - M N] measurable-vimage-Inr[of - M N]*
show *?case*
by (*simp add: indicator-vimage[symmetric] emeasure-copair-measure*)
next
case (*mult u c*)
then show *?case*
by(*simp add: measurable-copair-measure-iff nn-integral-cmult distrib-left*)
next
case (*add u v*)
then show *?case*
by(*simp add: nn-integral-add*)
next
case *h[measurable]:(seq U)*
have *inc:* $\bigwedge x. \text{incseq } (\lambda i. U i x)$
by (*metis h(3) incseq-def le-funE*)
have *lim:* $(\lambda i. U i x) \longrightarrow \text{Sup } (\text{range } U) x$ **for** *x*
by (*metis SUP-apply LIMSEQ-SUP[OF inc[of x]]*)
have $(\lambda i. (\int^+ x. U i x \partial(M \oplus_M N))) \longrightarrow (\int^+ x. (\text{Sup } (\text{range } U)) x \partial(M \oplus_M N))$
by(*intro nn-integral-LIMSEQ[OF - - lim]*) (*auto simp: h*)
moreover have $(\lambda i. (\int^+ x. U i x \partial(M \oplus_M N))) \longrightarrow (\int^+ x. \text{Sup } (\text{range } U) (\text{Inl } x) \partial M) + (\int^+ x. \text{Sup } (\text{range } U) (\text{Inr } x) \partial N)$
proof -
have $(\lambda i. (\int^+ x. U i x \partial(M \oplus_M N))) = (\lambda i. (\int^+ x. U i (\text{Inl } x) \partial M) + (\int^+ x. U i (\text{Inr } x) \partial N))$
by(*simp add: h*)
also have $\dots \longrightarrow (\int^+ x. \text{Sup } (\text{range } U) (\text{Inl } x) \partial M) + (\int^+ x. \text{Sup } (\text{range } U) (\text{Inr } x) \partial N)$
proof(*rule tendsto-add*)
have *inc:* $\bigwedge x. \text{incseq } (\lambda i. U i (\text{Inl } x))$
by (*metis h(3) incseq-def le-funE*)
have *lim:* $(\lambda i. U i (\text{Inl } x)) \longrightarrow \text{Sup } (\text{range } U) (\text{Inl } x)$ **for** *x*
by (*metis SUP-apply LIMSEQ-SUP[OF inc[of x]]*)
show $(\lambda i. (\int^+ x. U i (\text{Inl } x) \partial M)) \longrightarrow (\int^+ x. \text{Sup } (\text{range } U) (\text{Inl } x) \partial M)$
using *inc* **by**(*intro nn-integral-LIMSEQ[OF - - lim]*) (*auto simp: incseq-def intro!: le-funI*)
next
have *inc:* $\bigwedge x. \text{incseq } (\lambda i. U i (\text{Inr } x))$
by (*metis h(3) incseq-def le-funE*)
have *lim:* $(\lambda i. U i (\text{Inr } x)) \longrightarrow \text{Sup } (\text{range } U) (\text{Inr } x)$ **for** *x*

by (*metis SUP-apply LIMSEQ-SUP[OF inc[of x]]*)
show ($\lambda i. (\int^+ x. U i (Inr x) \partial N) \longrightarrow (\int^+ x. Sup (range U) (Inr x) \partial N)$)
using *inc* **by**(*intro nn-integral-LIMSEQ[OF - - lim]*) (*auto simp: incseq-def intro!: le-funI*)
qed
finally show *?thesis* .
qed
ultimately show *?case*
using *LIMSEQ-unique* **by** *blast*
qed

2.6 Integrability

lemma *integrable-copair-measure-iff*:

fixes $f :: 'a + 'b \Rightarrow 'c :: \{\text{banach, second-countable-topology}\}$
shows $integrable (M \oplus_M N) f \longleftrightarrow integrable M (\lambda x. f (Inl x)) \wedge integrable N (\lambda x. f (Inr x))$
by(*auto simp add: measurable-copair-measure-iff nn-integral-copair-measure integrable-iff-bounded*)

corollary *interable-copair-measureI*:

fixes $f :: 'a + 'b \Rightarrow 'c :: \{\text{banach, second-countable-topology}\}$
shows $integrable M (\lambda x. f (Inl x)) \implies integrable N (\lambda x. f (Inr x)) \implies integrable (M \oplus_M N) f$
by(*simp add: integrable-copair-measure-iff*)

2.7 The Lebesgue Integral

lemma *integral-copair-measure*:

fixes $f :: 'a + 'b \Rightarrow 'c :: \{\text{banach, second-countable-topology}\}$
assumes $integrable (M \oplus_M N) f$
shows $(\int x. f x \partial (M \oplus_M N)) = (\int x. f (Inl x) \partial M) + (\int x. f (Inr x) \partial N)$
using *assms*

proof *induction*

case $h[\text{measurable}]:(\text{base } A \ c)$
note $[\text{measurable}] = \text{measurable-vimage-Inl}[of \ - \ M \ N] \ \text{measurable-vimage-Inr}[of \ - \ M \ N]$
have $[\text{simp}]: integrable (M \oplus_M N) (\text{indicat-real } A) \ \text{integrable } M (\text{indicat-real } (Inl \ - \ A))$
 $\quad \text{integrable } N (\text{indicat-real } (Inr \ - \ A))$
using $h(2)$ **by**(*auto simp: emeasure-copair-measure*)
show *?case*
by(*cases c = 0*)
 $(\text{simp-all add: indicator-vimage[symmetric] \ measure-copair-measure \ measure-copair-measure[OF \ - \ h(2)] \ scaleR-left-distrib})$
next
case $(\text{add } f \ g)$
then show *?case*
by(*simp add: integrable-copair-measure-iff*)

next
case $ih:(\lim f s)$
have $(\lambda n. (\int x. s n x \partial(M \oplus_M N))) \longrightarrow (\int x. f x \partial(M \oplus_M N))$
using $ih(1-4)$ **by** $(\text{auto intro!}: \text{integral-dominated-convergence}[\text{where } w=\lambda x. 2 * \text{norm } (f x)])$
moreover have $(\lambda n. (\int x. s n x \partial(M \oplus_M N))) \longrightarrow (\int x. f (Inl x) \partial M) + (\int x. f (Inr x) \partial N)$
using $ih(1-4)$
by $(\text{auto intro!}: \text{integral-dominated-convergence}[\text{where } w=\lambda x. 2 * \text{norm } (f (Inl x))])$
 $\text{integral-dominated-convergence}[\text{where } w=\lambda x. 2 * \text{norm } (f (Inr x))]$ tendsto-add
 $\text{simp: } ih(5) \text{ integrable-copair-measure-iff measurable-copair-measure-iff borel-measurable-integrable space-copair-measure Inl InrI}$
ultimately show $?case$
using $LIMSEQ-unique$ **by** blast
qed

3 Coproduct Measures

definition $coPiM :: ['i \text{ set}, 'i \Rightarrow 'a \text{ measure}] \Rightarrow ('i \times 'a) \text{ measure}$ **where**
 $coPiM I Mi \equiv \text{measure-of}$
 $(SIGMA i:I. \text{space } (Mi i))$
 $\{A. A \subseteq (SIGMA i:I. \text{space } (Mi i)) \wedge (\forall i \in I. \text{Pair } i - ' A \in \text{sets } (Mi i))\}$
 $(\lambda A. (\sum_{\infty i \in I. \text{emeasure } (Mi i) (Pair i - ' A)))$

syntax

$-coPiM :: \text{pttrn} \Rightarrow 'i \text{ set} \Rightarrow 'a \text{ measure} \Rightarrow ('i \times 'a) \text{ measure} ((\exists \Pi_M - \in - / -) 10)$

translations

$\Pi_M x \in I. M \Leftrightarrow \text{CONST } coPiM I (\lambda x. M)$

3.1 The Measurable Space and Measurability

lemma

shows $\text{space-coPiM: space } (coPiM I Mi) = (SIGMA i:I. \text{space } (Mi i))$
and sets-coPiM:
 $\text{sets } (coPiM I Mi) = \text{sigma-sets } (SIGMA i:I. \text{space } (Mi i)) \{A. A \subseteq (SIGMA i:I. \text{space } (Mi i)) \wedge (\forall i \in I. \text{Pair } i - ' A \in \text{sets } (Mi i))\}$
and $\text{sets-coPiM-eq:sets } (coPiM I Mi) = \{A. A \subseteq (SIGMA i:I. \text{space } (Mi i)) \wedge (\forall i \in I. \text{Pair } i - ' A \in \text{sets } (Mi i))\}$
proof –
have $1:\{A. A \subseteq (SIGMA i:I. \text{space } (Mi i)) \wedge (\forall i \in I. \text{Pair } i - ' A \in \text{sets } (Mi i))\} \subseteq \text{Pow } (SIGMA i:I. \text{space } (Mi i))$
using $\text{sets.sets-into-space}$ **by** auto
show $\text{space } (coPiM I Mi) = (SIGMA i:I. \text{space } (Mi i))$
and $2:\text{sets } (coPiM I Mi)$
 $= \text{sigma-sets } (SIGMA i:I. \text{space } (Mi i)) \{A. A \subseteq (SIGMA i:I. \text{space } (Mi i)) \wedge (\forall i \in I. \text{Pair } i - ' A \in \text{sets } (Mi i))\}$

by(*auto simp: sets-measure-of[OF 1] space-measure-of[OF 1] coPiM-def*)
show $\text{sets } (\text{coPiM } I \text{ } Mi) = \{A. A \subseteq (\text{SIGMA } i:I. \text{space } (Mi \ i)) \wedge (\forall i \in I. \text{Pair } i - 'A \in \text{sets } (Mi \ i))\}$
proof –
have $\text{sigma-algebra } (\text{SIGMA } i:I. \text{space } (Mi \ i)) \{A. A \subseteq (\text{SIGMA } i:I. \text{space } (Mi \ i)) \wedge (\forall i \in I. \text{Pair } i - 'A \in \text{sets } (Mi \ i))\}$
proof(*subst Dynkin-system.sigma-algebra-eq-Int-stable*)
show $\text{Dynkin-system } (\text{SIGMA } i:I. \text{space } (Mi \ i)) \{A. A \subseteq (\text{SIGMA } i:I. \text{space } (Mi \ i)) \wedge (\forall i \in I. \text{Pair } i - 'A \in \text{sets } (Mi \ i))\}$
by *unfold-locales (auto simp: Pair-vimage-Sigma sets.Diff vimage-Diff vimage-Union 1)*
qed (*auto intro!: Int-stableI*)
thus *?thesis*
by(*auto simp: 2 intro!: sigma-algebra.sigma-sets-eq*)
qed
qed

lemma *sets-coPiM-cong*:

$I = J \implies (\bigwedge i. i \in I \implies \text{sets } (Mi \ i) = \text{sets } (Ni \ i)) \implies \text{sets } (\text{coPiM } I \text{ } Mi) = \text{sets } (\text{coPiM } J \text{ } Ni)$

by(*simp cong: sets-eq-imp-space-eq Sigma-cong add: sets-coPiM*)

lemma *measurable-coPiM2*:

assumes [*measurable*]: $\bigwedge i. i \in I \implies f \ i \in Mi \ i \rightarrow_M N$

shows $(\lambda(i,x). f \ i \ x) \in \text{coPiM } I \text{ } Mi \rightarrow_M N$

proof(*rule measurableI*)

fix A

assume [*measurable*]: $A \in \text{sets } N$

have [*simp*]:

$\bigwedge i. i \in I$

$\implies \text{Pair } i - '(\lambda(x, y). f \ x \ y) - 'A \cap \text{Pair } i - '(\text{SIGMA } i:I. \text{space } (Mi \ i)) = f \ i - 'A \cap \text{space } (Mi \ i)$

by *auto*

show $(\lambda(i, x). f \ i \ x) - 'A \cap \text{space } (\text{coPiM } I \text{ } Mi) \in \text{sets } (\text{coPiM } I \text{ } Mi)$

by(*auto simp: sets-coPiM space-coPiM*)

qed(*auto simp: space-coPiM measurable-space[OF assms]*)

lemma *measurable-Pair-coPiM[measurable (raw)]*:

assumes $i \in I$

shows $\text{Pair } i \in Mi \ i \rightarrow_M \text{coPiM } I \text{ } Mi$

proof(*rule measurable-sigma-sets*)

show $\{A. A \subseteq (\text{SIGMA } i:I. \text{space } (Mi \ i)) \wedge (\forall i \in I. \text{Pair } i - 'A \in \text{sets } (Mi \ i))\} \subseteq \text{Pow } (\text{SIGMA } i:I. \text{space } (Mi \ i))$

by *blast*

qed (*auto simp: assms sets-coPiM*)

lemma *measurable-Pair-coPiM'*:

assumes $i \in I (\lambda(i,x). f \ i \ x) \in \text{coPiM } I \text{ } Mi \rightarrow_M N$

shows $f \ i \in Mi \ i \rightarrow_M N$

using *measurable-compose*[*OF measurable-Pair-coPiM assms(2)*] *assms(1)* **by**
fast

lemma *measurable-copair-iff*: $(\lambda(i,x). f\ i\ x) \in \text{coPiM } I\ Mi \rightarrow_M N \iff (\forall i \in I. f\ i \in Mi\ i \rightarrow_M N)$

by(*auto intro!*: *measurable-coPiM2 simp: measurable-Pair-coPiM'*)

lemma *measurable-copair-iff'*: $f \in \text{coPiM } I\ Mi \rightarrow_M N \iff (\forall i \in I. (\lambda x. f\ (i, x)) \in Mi\ i \rightarrow_M N)$

using *measurable-copair-iff*[*of curry f*] **by**(*simp add: split-beta' curry-def*)

lemma *coPair-inverse-space-unit*:

$i \in I \implies A \in \text{sets } (\text{coPiM } I\ Mi) \implies \text{Pair } i - ' A \cap \text{space } (Mi\ i) = \text{Pair } i - ' A$
using *sets.sets-into-space* **by**(*fastforce simp: space-coPiM*)

lemma *measurable-Pair-vimage*:

assumes $i \in I\ A \in \text{sets } (\text{coPiM } I\ Mi)$

shows $\text{Pair } i - ' A \in \text{sets } (Mi\ i)$

using *measurable-sets*[*OF measurable-Pair-coPiM*][*OF assms(1)*] *assms(2)*

by (*simp add: coPair-inverse-space-unit*[*OF assms*])

lemma *measurable-Sigma-singleton*[*measurable (raw)*]:

$\bigwedge i\ A. i \in I \implies A \in \text{sets } (Mi\ i) \implies \{i\} \times A \in \text{sets } (\text{coPiM } I\ Mi)$

using *sets.sets-into-space sets-coPiM* **by** *fastforce*

lemma *sets-coPiM-countable*:

assumes *countable I*

shows $\text{sets } (\text{coPiM } I\ Mi) = \text{sigma-sets } (\text{SIGMA } i:I. \text{space } (Mi\ i)) (\bigcup i \in I. (\times) \{i\} - ' (\text{sets } (Mi\ i)))$

unfolding *sets-coPiM*

proof(*safe intro!*: *sigma-sets-eqI*)

fix *a*

assume $h:a \subseteq (\text{SIGMA } i:I. \text{space } (Mi\ i)) \forall i \in I. \text{Pair } i - ' a \in \text{sets } (Mi\ i)$

then have $a = (\bigcup i \in I. \{i\} \times \text{Pair } i - ' a)$

by *auto*

moreover have $(\bigcup i \in I. \{i\} \times \text{Pair } i - ' a) \in \text{sigma-sets } (\text{SIGMA } i:I. \text{space } (Mi\ i)) (\bigcup i \in I. (\times) \{i\} - ' (\text{sets } (Mi\ i)))$

using *h(2)* **by**(*auto intro!*: *sigma-sets-UNION*[*OF countable-image*][*OF assms*])

ultimately show $a \in \text{sigma-sets } (\text{SIGMA } i:I. \text{space } (Mi\ i)) (\bigcup i \in I. (\times) \{i\} - ' (\text{sets } (Mi\ i)))$

by *argo*

qed(*use sets.sets-into-space in fastforce*)

lemma *measurable-coPiM1'*:

assumes *countable I*

and [*measurable*]: $a \in N \rightarrow_M \text{count-space } I \bigwedge i. i \in a - ' (\text{space } N) \implies g\ i \in N \rightarrow_M Mi\ i$

shows $(\lambda x. (a\ x, g\ (a\ x)\ x)) \in N \rightarrow_M \text{coPiM } I\ Mi$

proof(*safe intro!*: *measurable-sigma-sets*[*OF sets-coPiM-countable*][*OF assms(1)*])

```

fix i B
assume iB[measurable]:i ∈ I B ∈ sets (Mi i)
show (λx. (a x, g (a x) x)) -‘ ({i} × B) ∩ space N ∈ sets N
proof(cases i ∈ a -‘ (space N))
  assume [measurable]:i ∈ a -‘ (space N)
  have (λx. (a x, g (a x) x)) -‘ ({i} × B) ∩ space N = (a -‘ {i} ∩ space N) ∩
(g i -‘ B ∩ space N)
    by auto
  also have ... ∈ sets N
    by simp
  finally show ?thesis .
next
assume i ∉ a -‘ (space N)
then have (λx. (a x, g (a x) x)) -‘ ({i} × B) ∩ space N = {}
  using measurable-space[OF assms(2)] by blast
thus ?thesis
  by simp
qed
qed(use measurable-space[OF assms(2)] measurable-space[OF assms(3)] sets.sets-into-space
in fastforce)+

```

```

lemma measurable-coPiM1:
assumes countable I
  and a ∈ N →M count-space I ∧ i. i ∈ I ⇒ g i ∈ N →M Mi i
shows (λx. (a x, g (a x) x)) ∈ N →M coPiM I Mi
using measurable-space[OF assms(2)] by(auto intro!: measurable-coPiM1' assms)

```

```

lemma measurable-coPiM1-elements:
assumes countable I and [measurable]:f ∈ N →M coPiM I Mi
obtains a g
  where a ∈ N →M count-space I
    ∧ i. i ∈ I ⇒ space (Mi i) ≠ {} ⇒ g i ∈ N →M Mi i
    f = (λx. (a x, g (a x) x))
proof -
define a where a ≡ fst ∘ f
have 1[measurable]:a ∈ N →M count-space I
proof(safe intro!: measurable-count-space-eq-countable[THEN iffD2] assms)
  fix i
  assume i:i ∈ I
  have a -‘ {i} ∩ space N = f -‘ ({i} × space (Mi i)) ∩ space N
    using measurable-space[OF assms(2)] by(fastforce simp: a-def space-coPiM)
  also have ... ∈ sets N
    using i by auto
  finally show a -‘ {i} ∩ space N ∈ sets N .
next
show ∧x. x ∈ space N ⇒ a x ∈ I
  using measurable-space[OF assms(2)] by(fastforce simp: space-coPiM a-def)
qed
define g where g ≡ (λi x. if a x = i then snd (f x) else (SOME y. y ∈ space

```

```

(Mi i))
  have 2:g i ∈ N →M Mi i if i:i ∈ I and ne:space (Mi i) ≠ {} for i
    unfolding g-def
  proof(safe intro!: measurable-If-restrict-space-iff[THEN iffD2] measurable-const
some-in-eq[THEN iffD2] ne)
    show (λx. snd (f x)) ∈ restrict-space N {x. a x = i} →M Mi i
    proof(safe intro!: measurableI)
      show ∧x. x ∈ space (restrict-space N {x. a x = i}) ⇒ snd (f x) ∈ space
(Mi i)
      using measurable-space[OF assms(2)] by(fastforce simp: space-restrict-space
a-def space-coPiM)
    next
    fix A
    assume [measurable]:A ∈ sets (Mi i)
    have (λx. snd (f x)) -‘ A ∩ space (restrict-space N {x. a x = i}) = f -‘ ({i}
× A) ∩ space N
    using i measurable-space[OF assms(2)] by(fastforce simp: space-restrict-space
a-def space-coPiM)
    also have ... ∈ sets N
    using i by simp
    finally show (λx. snd (f x)) -‘ A ∩ space (restrict-space N {x. a x = i})
      ∈ sets (restrict-space N {x. a x = i})
      by(auto simp: sets-restrict-space space-restrict-space)
    qed
  qed(use i ne in auto)
  have 3:f = (λx. (a x, g (a x) x))
    by(auto simp: a-def g-def)
  show ?thesis
    using 1 2 3 that by blast
qed

```

3.2 Measures

lemma *emeasure-coPiM*:

```

  assumes A ∈ sets (coPiM I Mi)
  shows emeasure (coPiM I Mi) A = (∑∞ i∈I. emeasure (Mi i) (Pair i -‘ A))
proof(rule emeasure-measure-of)
  show {A. A ⊆ (SIGMA i:I. space (Mi i)) ∧ (∀ i∈I. Pair i -‘ A ∈ sets (Mi i))}
  ⊆ Pow (SIGMA i:I. space (Mi i))
    by blast
  next
  note measurable-Pair-vimage[of - I - Mi,measurable (raw)]
  show countably-additive (sets (coPiM I Mi)) (λa. ∑∞ i∈I. emeasure (Mi i) (Pair
i -‘ a))
    unfolding countably-additive-def
  proof safe
    fix A :: nat ⇒ -
    assume A:range A ⊆ sets (coPiM I Mi) disjoint-family A
    then have [measurable]:∧n. A n ∈ sets (coPiM I Mi)

```


by *blast*
show $(\sum n. \sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (Pair\ i - ' A\ n))$
 $= (\sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (Pair\ i - ' \bigcup (\text{range } A)))$ (**is** *?lhs = ?rhs*)
proof –
 have *?lhs* $= (\sum_{\infty n \in UNIV}. \sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (Pair\ i - ' A\ n))$
 by(*auto intro!*: *infsum-eq-suminf[symmetric] nonneg-summable-on-complete*)
 also have ... $= (\sum_{\infty i \in I}. \sum_{\infty n \in UNIV}. \text{emeasure } (Mi\ i) (Pair\ i - ' A\ n))$
 by(*rule infsum-swap-ennreal*)
 also have ... $= ?rhs$
 proof(*rule infsum-cong*)
 fix *i*
 assume $i \in I$
 then have $(\sum n. Mi\ i (Pair\ i - ' A\ n)) = Mi\ i (\bigcup n. Pair\ i - ' A\ n)$
 using *A(2)* by(*intro suminf-emeasure*) (*auto simp: disjoint-family-on-def*)
 also have ... $= Mi\ i (Pair\ i - ' \bigcup (\text{range } A))$
 by (*metis vimage-UN*)
 finally show $(\sum_{\infty n}. \text{emeasure } (Mi\ i) (Pair\ i - ' A\ n)) = \text{emeasure } (Mi\ i)$
 (*Pair\ i - ' \bigcup (\text{range } A)*)
 by(*auto simp: infsum-eq-suminf[OF nonneg-summable-on-complete]*)
 qed
 finally show *?thesis* .
 qed
qed
next
 show $A \in \text{sets } (coPiM\ I\ Mi)$
 by *fact*
qed(*auto simp: positive-def coPiM-def*)

corollary *emeasure-coPiM-space*:
 $\text{emeasure } (coPiM\ I\ Mi) (\text{space } (coPiM\ I\ Mi)) = (\sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (\text{space } (Mi\ i)))$
proof –
 have [*simp*]: $\bigwedge i. i \in I \implies Pair\ i - ' \text{space } (coPiM\ I\ Mi) = \text{space } (Mi\ i)$
 by(*auto simp: space-coPiM*)
 show *?thesis*
 by(*auto simp: emeasure-coPiM intro!: infsum-cong*)
qed

lemma *emeasure-coPiM-coproj*:
 assumes [*measurable*]: $i \in I\ A \in \text{sets } (Mi\ i)$
 shows $\text{emeasure } (coPiM\ I\ Mi) (\{i\} \times A) = \text{emeasure } (Mi\ i)\ A$
proof –
 have $\text{emeasure } (coPiM\ I\ Mi) (\{i\} \times A) = (\sum_{\infty j \in I}. \text{emeasure } (Mi\ j) (\text{if } j = i \text{ then } A \text{ else } \{\}))$
 by(*simp add: emeasure-coPiM*)
 also have ... $= (\sum_{\infty j \in (I - \{i\}) \cup \{i\}}. \text{emeasure } (Mi\ j) (\text{if } j = i \text{ then } A \text{ else } \{\}))$
 by(*rule arg-cong[where f=infsum -]*) (*use assms in auto*)
 also have ... $= (\sum_{\infty j \in I - \{i\}}. \text{emeasure } (Mi\ j) (\text{if } j = i \text{ then } A \text{ else } \{\}))$

$+ (\sum_{\infty j \in \{i\}}. \text{emeasure } (Mi\ j) \text{ (if } j = i \text{ then } A \text{ else } \{\}))$
by(rule *infsum-Un-disjoint*) (auto intro!: *nonneg-summable-on-complete*)
also have ... = *emeasure* (Mi i) A
proof –
have $(\sum_{\infty j \in I - \{i\}}. \text{emeasure } (Mi\ j) \text{ (if } j = i \text{ then } A \text{ else } \{\})) = 0$
by (rule *infsum-0*) *simp*
thus ?thesis **by** *simp*
qed
finally show ?thesis .
qed

lemma *measure-coPiM-coproj*: $i \in I \implies A \in \text{sets } (Mi\ i) \implies \text{measure } (coPiM\ I\ Mi) (\{i\} \times A) = \text{measure } (Mi\ i) A$
by(*simp add: emeasure-coPiM-coproj measure-def*)

lemma *emeasure-coPiM-less-top-summable*:

assumes [*measurable*]: $A \in \text{sets } (coPiM\ I\ Mi)$ *emeasure* (coPiM I Mi) A < ∞
shows $(\lambda i. \text{measure } (Mi\ i) (Pair\ i - 'A))$ *summable-on* I
proof –
have *: $(\sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (Pair\ i - 'A)) < \text{top}$
using *assms(2)* **by**(*simp add: emeasure-coPiM*)
from *infsum-less-top-dest[OF this]* **have** *ifin*: $\bigwedge i. i \in I \implies \text{emeasure } (Mi\ i) (Pair\ i - 'A) < \text{top}$
by *simp*
with * **have** $(\sum_{\infty i \in I}. \text{ennreal } (\text{measure } (Mi\ i) (Pair\ i - 'A))) < \text{top}$
by (*metis* (*mono-tags*, *lifting*) *emeasure-eq-ennreal-measure infsum-cong top.not-eq-extremum*)
thus ?thesis
by(auto intro!: *bounded-infsum-summable*)
qed

lemma *measure-coPiM*:

assumes [*measurable*]: $A \in \text{sets } (coPiM\ I\ Mi)$ *emeasure* (coPiM I Mi) A < ∞
shows $\text{measure } (coPiM\ I\ Mi) A = (\sum_{\infty i \in I}. \text{measure } (Mi\ i) (Pair\ i - 'A))$
proof(*subst ennreal-inj[symmetric]*)
have *: $(\sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (Pair\ i - 'A)) < \text{top}$
using *assms(2)* **by**(*simp add: emeasure-coPiM*)
from *infsum-less-top-dest[OF this]* **have** *ifin*: $\bigwedge i. i \in I \implies \text{emeasure } (Mi\ i) (Pair\ i - 'A) < \text{top}$
by *simp*
show $\text{ennreal } (\text{measure } (coPiM\ I\ Mi) A) = \text{ennreal } (\sum_{\infty i \in I}. \text{measure } (Mi\ i) (Pair\ i - 'A))$ (*is* ?lhs = ?rhs)
proof –
have ?lhs = *emeasure* (coPiM I Mi) A
using *assms* **by**(auto intro!: *emeasure-eq-ennreal-measure[symmetric]*)
also have ... = $(\sum_{\infty i \in I}. \text{emeasure } (Mi\ i) (Pair\ i - 'A))$
by(*simp add: emeasure-coPiM*)
also have ... = $(\sum_{\infty i \in I}. \text{ennreal } (\text{measure } (Mi\ i) (Pair\ i - 'A)))$
using *ifin* **by**(*fastforce intro!: infsum-cong emeasure-eq-ennreal-measure*)
also have ... = ?rhs

```

    by(simp add: infsum-ennreal-eq[OF emeasure-coPiM-less-top-summable[OF
assms]])
    finally show ?thesis .
  qed
qed(auto intro!: infsum-nonneg)

```

3.3 Non-Negative Integral

```

lemma nn-integral-coPiM:
  assumes f ∈ borel-measurable (coPiM I Mi)
  shows (∫+x. f x ∂coPiM I Mi) = (∑∞i∈I. (∫+x. f (i, x) ∂Mi i))
  using assms
proof induction
  case (cong f g)
  moreover hence ∧i x. i ∈ I ⇒ x ∈ space (Mi i) ⇒ f (i, x) = g (i, x)
  by(auto simp: space-coPiM)
  ultimately show ?case
  by(simp cong: nn-integral-cong infsum-cong)
next
  case [measurable]:(set A)
  note [measurable] = measurable-Pair-vimage[OF - this]
  show ?case
  by(simp add: indicator-vimage[symmetric] emeasure-coPiM cong: infsum-cong)
next
  case (add u v)
  then show ?case
  by(simp add: nn-integral-add infsum-add nonneg-summable-on-complete cong:
infsum-cong)
next
  case (mult u c)
  then show ?case
  by(simp add: nn-integral-cmult infsum-cmult-right-ennreal cong: infsum-cong)
next
  case ih[measurable]:(seq U)
  show ?case (is ?lhs = ?rhs)
  proof -
    have ?lhs = (∫+x. (SUP j. U j x) ∂coPiM I Mi)
    by(auto intro!: nn-integral-cong simp: SUP-apply[symmetric])
    also have ... = (SUP j. (∫+x. U j x ∂coPiM I Mi))
    by(auto intro!: nn-integral-monotone-convergence-SUP ih(?))
    also have ... = (SUP j. (∑∞i∈I. (∫+x. U j (i, x) ∂Mi i)))
    by(simp add: ih)
    also have ... = (∑∞i∈I. (SUP j. (∫+x. U j (i, x) ∂Mi i)))
    using ih(?)) by(auto intro!: ennreal-infsum-Sup-eq[symmetric] incseq-nn-integral
simp: mono-compose)
    also have ... = (∑∞i∈I. (∫+x. (SUP j. U j (i, x)) ∂Mi i))
    using ih(?)) by(auto intro!: infsum-cong nn-integral-monotone-convergence-SUP[symmetric]
mono-compose)
    also have ... = ?rhs

```

```

    by(simp add: SUP-apply[symmetric])
  finally show ?thesis .
qed
qed

```

3.4 Integrability

lemma

```

  fixes f :: - => 'b::{banach, second-countable-topology}
  assumes integrable (coPiM I Mi) f
  shows integrable-coPiM-dest-sum:( $\sum_{\infty i \in I. (\int^+ x. \text{norm } (f (i, x)) \partial Mi i) < \infty$ )
  and integrable-coPiM-dest-integrable:  $\bigwedge i. i \in I \implies \text{integrable } (Mi i) (\lambda x. f (i, x))$ 
  and integrable-coPiM-summable-norm:  $(\lambda i. (\int x. \text{norm } (f (i, x)) \partial Mi i))$  summable-on I
  and integrable-coPiM-abs-summable: Infinite-Sum.abs-summable-on  $(\lambda i. (\int x. f (i, x) \partial Mi i)) I$ 
  and integrable-coPiM-summable:  $(\lambda i. (\int x. f (i, x) \partial Mi i))$  summable-on I
proof -
  show fin:( $\sum_{\infty i \in I. (\int^+ x. \text{norm } (f (i, x)) \partial Mi i) < \infty$ )
    using assms by(auto simp: integrable-iff-bounded nn-integral-coPiM)
  thus integ: $i \in I \implies \text{integrable } (Mi i) (\lambda x. f (i, x))$  for i
    using assms by(auto simp: integrable-iff-bounded intro!: infsum-less-top-dest[of - - i])
  show summable:( $\lambda i. (\int x. \text{norm } (f (i, x)) \partial Mi i)$ ) summable-on I
    using nn-integral-eq-integral[OF integrable-norm[OF integ]] fin
    by(auto intro!: bounded-infsum-summable cong: infsum-cong)
  show Infinite-Sum.abs-summable-on  $(\lambda i. (\int x. f (i, x) \partial Mi i)) I$ 
    by(rule summable-on-comparison-test[OF summable]) auto
  thus  $(\lambda i. (\int x. f (i, x) \partial Mi i))$  summable-on I
    using abs-summable-summable by fastforce
qed

```

3.5 The Lebesgue Integral

lemma *integral-coPiM*:

```

  fixes f :: - => 'b::{banach, second-countable-topology}
  assumes integrable (coPiM I Mi) f
  shows  $(\int x. f x \partial coPiM I Mi) = (\sum_{\infty i \in I. (\int x. f (i, x) \partial Mi i))$ 
  using assms
proof induction
  case h[measurable]:(base A c)
  note [measurable] = measurable-Pair-vimage[OF - this(1)]
  have [simp]: integrable (coPiM I Mi) (indicat-real A)
     $\bigwedge i. i \in I \implies \text{integrable } (Mi i) (\text{indicat-real } (Pair i - ' A))$ 
    using h(2) by(auto simp: emeasure-coPiM dest: infsum-less-top-dest)
  show ?case
    using h(2) emeasure-coPiM-less-top-summable[OF h]
    by(cases c = 0)

```

```

      (auto simp: measure-coPiM indicator-vimage[symmetric] infsum-scaleR-left[symmetric]
cong: infsum-cong)
next
case h:(add f g)
show ?case (is ?lhs = ?rhs)
proof -
  have ?lhs = (∑∞ i∈I. (∫ x. f (i, x) ∂Mi i)) + (∑∞ i∈I. (∫ x. g (i, x) ∂Mi i))
  using h by simp
  also have ... = (∑∞ i∈I. (∫ x. f (i, x) ∂Mi i) + (∫ x. g (i, x) ∂Mi i))
  by(auto intro!: infsum-add[symmetric] integrable-coPiM-summable h)
  also have ... = ?rhs
  using h
  by(auto intro!: infsum-cong Bochner-Integration.integral-add[symmetric] inte-
grable-coPiM-dest-integrable)
  finally show ?thesis .
qed
next
case ih:(lim f fn)
note [measurable,simp] = ih(1-4)
show ?case (is ?lhs = ?rhs)
proof -
  have ?lhs = lim (λn. (∫ x. fn n x ∂(coPiM I Mi)))
  by(auto intro!: limI[symmetric] integral-dominated-convergence[where w=λx.
2 * norm (f x)])
  also have ... = lim (λn. (∑∞ i∈I. (∫ x. fn n (i, x) ∂Mi i)))
  by(simp add: ih(5))
  also have ... = lim (λn. (∫ i. (∫ x. fn n (i, x) ∂Mi i) ∂count-space I))
  by(simp add: integrable-coPiM-abs-summable infsum-eq-integral)
  also have ... = (∫ i. (∫ x. f (i, x) ∂Mi i) ∂count-space I)
  proof(intro limI integral-dominated-convergence[where w=λi. (∫ x. 2 * norm
(f (i, x)) ∂Mi i)] AE-I2 )
    show integrable (count-space I) (λi. (∫ x. 2 * norm (f (i, x)) ∂Mi i))
    by(auto simp: abs-summable-on-integrable-iff[symmetric] integrable-coPiM-summable-norm[OF
ih(4)])
  next
  show i ∈ space (count-space I) ⇒ (λn. (∫ x. fn n (i, x) ∂Mi i)) → (∫ x.
f (i, x) ∂Mi i) for i
  by(auto intro!: integral-dominated-convergence[where w=λx. 2*norm (f (i,
x))]] integrable-coPiM-dest-integrable
simp: space-coPiM)
  next
  show i ∈ space (count-space I) ⇒ norm ((∫ x. fn n (i, x) ∂Mi i)) ≤ (∫ x.
2 * norm (f (i, x)) ∂Mi i) for n i
  by(rule order.trans[where b=(∫ x. norm (fn n (i, x)) ∂Mi i)])
  (auto simp: space-coPiM
simp del: Bochner-Integration.integral-mult-right-zero Bochner-Integration.integral-mult-right
intro!: integral-mono integrable-coPiM-dest-integrable)
qed simp-all
also have ... = ?rhs

```

by(*simp add: infsum-eq-integral integrable-coPiM-abs-summable*)
 finally show *?thesis* .
 qed
 qed

3.6 Finite Coproduct Measures

lemma *emeasure-coPiM-finite*:

assumes *finite I A ∈ sets (coPiM I Mi)*
 shows *emeasure (coPiM I Mi) A = (∑ i∈I. emeasure (Mi i) (Pair i -‘ A))*
 using *assms* by(*simp add: emeasure-coPiM*)

lemma *emeasure-coPiM-finite-space*:

finite I ⇒ emeasure (coPiM I Mi) (space (coPiM I Mi)) = (∑ i∈I. emeasure (Mi i) (space (Mi i)))
 by(*simp add: emeasure-coPiM-space*)

lemma *measure-coPiM-finite*:

assumes *finite I A ∈ sets (coPiM I Mi) emeasure (coPiM I Mi) A < ∞*
 shows *measure (coPiM I Mi) A = (∑ i∈I. measure (Mi i) (Pair i -‘ A))*
 using *assms(3)* by(*simp add: emeasure-coPiM-finite[OF assms(1,2)] measure-def enn2real-sum assms(1)*)

lemma *nn-integral-coPiM-finite*:

assumes *finite I f ∈ borel-measurable (coPiM I Mi)*
 shows *(∫⁺x. f x ∂(coPiM I Mi)) = (∑ i∈I. (∫⁺x. f (i, x) ∂(Mi i)))*
 by(*simp add: nn-integral-coPiM assms*)

lemma *integrable-coPiM-finite-iff*:

fixes *f :: - ⇒ ‘c::{banach, second-countable-topology}*
 shows *finite I ⇒ integrable (coPiM I Mi) f ↔ (∀ i∈I. integrable (Mi i) (λx. f (i, x)))*
 using *measurable-copair-iff'[of f I Mi borel]*
 by(*auto simp: integrable-iff-bounded nn-integral-coPiM-finite*)

lemma *integral-coPiM-finite*:

fixes *f :: - ⇒ ‘c::{banach, second-countable-topology}*
 assumes *finite I integrable (coPiM I Mi) f*
 shows *(∫ x. f x ∂(coPiM I Mi)) = (∑ i∈I. (∫ x. f (i, x) ∂(Mi i)))*
 by(*auto simp: assms integral-coPiM*)

3.7 Countable Infinite Coproduct Measures

lemma *emeasure-coPiM-countable-infinite*:

assumes [*measurable*]: *bij-betw from-n (UNIV :: nat set) I A ∈ sets (coPiM I Mi)*
 shows *emeasure (coPiM I Mi) A = (∑ n. emeasure (Mi (from-n n)) (Pair (from-n n) -‘ A))*
proof –
 have *I:countable I*

using *assms(1) countableI-bij* **by** *blast*
have [*measurable,simp*]:*Pair (from-n n) - ' A ∈ sets (Mi (from-n n)) from-n n*
 $\in I$ **for** *n*
using *bij-betwE[OF assms(1)]* **by**(*auto intro!*: *measurable-Pair-vimage*[**where**
 $I=I$])
have *emeasure (coPiM I Mi) A = emeasure (coPiM I Mi) ($\bigcup n. \{from-n n\} \times$*
Pair (from-n n) - ' A)
using *sets.sets-into-space[OF assms(2)] assms(1)*
by(*fastforce intro!*: *arg-cong*[**where** $f=emeasure (coPiM I Mi)$
simp: space-coPiM bij-betw-def])
also have $\dots = (\sum n. emeasure (Mi (from-n n)) (Pair (from-n n) - ' A))$
using *injD[OF bij-betw-imp-inj-on[OF assms(1)]]*
by(*subst suminf-emeasure[symmetric]*)
(auto simp: disjoint-family-on-def emeasure-coPiM-coproj intro!: suminf-cong)
finally show *?thesis .*
qed

lemmas *emeasure-coPiM-countable-infinite' = emeasure-coPiM-countable-infinite*[*OF*
bij-betw-from-nat-into]
lemmas *emeasure-coPiM-nat = emeasure-coPiM-countable-infinite*[*OF* *bij-id,simplified*]

lemma *measure-coPiM-countable-infinite:*

assumes [*measurable,simp*]: *bij-betw from-n (UNIV :: nat set) I A ∈ sets (coPiM*
 $I Mi)$
and *emeasure (coPiM I Mi) A < ∞*
shows *measure (coPiM I Mi) A = ($\sum n. measure (Mi (from-n n)) (Pair (from-n$*
 $n) - ' A)$) (**is** *?lhs = ?rhs*)
and *summable ($\lambda n. measure (Mi (from-n n)) (Pair (from-n n) - ' A)$)*
proof –
have *ennreal ?lhs = emeasure (coPiM I Mi) A*
using *assms(3) by(auto intro!: emeasure-eq-ennreal-measure[symmetric])*
also have $\dots = (\sum n. emeasure (Mi (from-n n)) (Pair (from-n n) - ' A))$
by(*simp add: emeasure-coPiM-countable-infinite*)
also have $\dots = (\sum n. ennreal (measure (Mi (from-n n)) (Pair (from-n n) - ' A)))$
using *assms(3) ennreal-suminf-lessD top.not-eq-extremum*
by(*auto intro!: suminf-cong emeasure-eq-ennreal-measure*
simp: emeasure-coPiM-countable-infinite[OF assms(1)]])
finally have $*:ennreal ?lhs = (\sum n. ennreal (measure (Mi (from-n n)) (Pair$
 $(from-n n) - ' A))) .$
thus $**[simp]: summable (\lambda n. measure (Mi (from-n n)) (Pair (from-n n) - ' A))$
by(*auto intro!: summable-suminf-not-top*)
show *?lhs = ?rhs*
proof(*subst ennreal-inj[symmetric]*)
have *ennreal ?lhs = ($\sum n. ennreal (measure (Mi (from-n n)) (Pair (from-n n)$*
 $- ' A)))$
by *fact*
also have $\dots = ennreal ?rhs$
using *assms(3) by(auto intro!: suminf-ennreal2)*

finally show $ennreal ?lhs = ennreal ?rhs$.
qed(*simp-all add: suminf-nonneg*)
qed

lemmas *measure-coPiM-countable-infinite'* = *measure-coPiM-countable-infinite*[*OF bij-betw-from-nat-into*]
lemmas *measure-coPiM-nat* = *measure-coPiM-countable-infinite*[*OF bij-id,simplified id-apply*]

lemma *nn-integral-coPiM-countable-infinite*:
assumes [*measurable*]:*bij-betw from-n (UNIV :: nat set) I f* \in *borel-measurable (coPiM I Mi)*
shows $(\int^{+x}. f x \partial(\text{coPiM } I \text{ } Mi)) = (\sum n. (\int^{+x}. f (\text{from-n } n, x) \partial(Mi (\text{from-n } n))))$ (*is - = ?rhs*)
proof -
have $(\int^{+x}. f x \partial(\text{coPiM } I \text{ } Mi)) = (\sum_{\infty} i \in I. (\int^{+x}. f (i, x) \partial Mi i))$
by(*simp add: nn-integral-coPiM*)
also have $\dots = (\sum_{\infty} i \in \text{from-n } 'UNIV. (\int^{+x}. f (i, x) \partial Mi i))$
by(*rule arg-cong[where f=infsum -]*) (*metis assms(1) bij-betw-def*)
also have $\dots = (\sum_{\infty} n \in UNIV. (\int^{+x}. f (\text{from-n } n, x) \partial(Mi (\text{from-n } n))))$
by(*rule infsum-reindex[simplified comp-def]*) (*use assms(1) bij-betw-imp-inj-on*)
in blast
also have $\dots = ?rhs$
by(*auto intro!: infsum-eq-suminf nonneg-summable-on-complete*)
finally show *?thesis* .

qed
lemmas *nn-integral-coPiM-countable-infinite'* = *nn-integral-coPiM-countable-infinite*[*OF bij-betw-from-nat-into*]
lemmas *nn-integral-coPiM-nat* = *nn-integral-coPiM-countable-infinite*[*OF bij-id,simplified*]

lemma
fixes $f :: - \Rightarrow 'b :: \{banach, second-countable-topology\}$
assumes *bij-betw from-n (UNIV :: nat set) I integrable (coPiM I Mi) f*
shows *integrable-coPiM-countable-infinite-dest-sum*: $(\sum n. (\int^{+x}. norm (f (\text{from-n } n, x)) \partial(Mi (\text{from-n } n)))) < \infty$
and *integrable-coPiM-countable-infinite-dest'*: $\bigwedge n. \text{integrable } (Mi (\text{from-n } n)) (\lambda x. f (\text{from-n } n, x))$
using *ennreal-suminf-lessD assms(1,2) bij-betwE[OF assms(1)]*
by(*auto simp: integrable-iff-bounded nn-integral-coPiM-countable-infinite*)

lemmas *integrable-coPiM-countable-infinite-dest-sum'* = *integrable-coPiM-countable-infinite-dest-sum*[*OF bij-betw-from-nat-into*]
lemmas *integrable-coPiM-countable-infinite-dest''* = *integrable-coPiM-countable-infinite-dest'*[*OF bij-betw-from-nat-into*]
lemmas *integrable-coPiM-nat-dest-sum* = *integrable-coPiM-countable-infinite-dest-sum*[*OF bij-id,simplified id-apply*]
lemmas *integrable-coPiM-nat-dest* = *integrable-coPiM-countable-infinite-dest'*[*OF bij-id,simplified id-apply*]

lemma

fixes $f :: - \Rightarrow 'b::\{\text{banach, second-countable-topology}\}$
assumes $\text{bij-betw from-n (UNIV :: nat set) I integrable (coPiM I Mi) f}$
shows $\text{integrable-coPiM-countable-infinite-summable-norm: summable } (\lambda n. (\int x. \text{norm } (f \text{ (from-n } n, x)) \partial(\text{Mi (from-n } n))))$
and $\text{integrable-coPiM-countable-infinite-summable-norm': summable } (\lambda n. \text{norm } (\int x. f \text{ (from-n } n, x) \partial(\text{Mi (from-n } n))))$
and $\text{integrable-coPiM-countable-infinite-summable: summable } (\lambda n. (\int x. f \text{ (from-n } n, x) \partial(\text{Mi (from-n } n))))$

proof –

show $*:\text{summable } (\lambda n. (\int x. \text{norm } (f \text{ (from-n } n, x)) \partial(\text{Mi (from-n } n))))$
using $\text{integrable-coPiM-countable-infinite-dest-sum[OF assms]}$
 $\text{nn-integral-eq-integral[OF integrable-norm[OF integrable-coPiM-countable-infinite-dest'[OF assms]]]}$
by $(\text{auto intro!: summable-suminf-not-top})$
show $\text{summable } (\lambda n. \text{norm } (\int x. f \text{ (from-n } n, x) \partial(\text{Mi (from-n } n))))$
by $(\text{rule summable-comparison-test-ev[OF - *] auto})$
thus $\text{summable } (\lambda n. (\int x. f \text{ (from-n } n, x) \partial(\text{Mi (from-n } n))))$
using $\text{summable-norm-cancel by force}$

qed

lemmas $\text{integrable-coPiM-countable-infinite-summable-norm''}$

$= \text{integrable-coPiM-countable-infinite-summable-norm[OF bij-betw-from-nat-into]}$

lemmas $\text{integrable-coPiM-countable-infinite-summable-norm'''}$

$= \text{integrable-coPiM-countable-infinite-summable-norm'[OF bij-betw-from-nat-into]}$

lemmas $\text{integrable-coPiM-countable-infinite-summable}'$

$= \text{integrable-coPiM-countable-infinite-summable[OF bij-betw-from-nat-into]}$

lemmas $\text{integrable-coPiM-nat-summable-norm}$

$= \text{integrable-coPiM-countable-infinite-summable-norm[OF bij-id,simplified id-apply]}$

lemmas $\text{integrable-coPiM-nat-summable-norm}'$

$= \text{integrable-coPiM-countable-infinite-summable-norm'[OF bij-id,simplified id-apply]}$

lemmas $\text{integrable-coPiM-nat-summable}$

$= \text{integrable-coPiM-countable-infinite-summable[OF bij-id,simplified id-apply]}$

lemma

fixes $f :: - \Rightarrow 'b::\{\text{banach, second-countable-topology}\}$
assumes $\text{countable I infinite I integrable (coPiM I Mi) f}$
shows $\text{integrable-coPiM-countable-infinite-dest: } \bigwedge i. i \in I \implies \text{integrable } (Mi \ i)$
 $(\lambda x. f \ (i, x))$
using $\text{integrable-coPiM-countable-infinite-dest'[OF bij-betw-from-nat-into[OF assms(1,2)]]}$
 assms(3)
by $(\text{meson assms(1) countable-all})$

lemma $\text{integrable-coPiM-countable-infiniteI:}$

fixes $f :: - \Rightarrow 'b::\{\text{banach, second-countable-topology}\}$
assumes $\text{bij-betw from-n (UNIV :: nat set) I } \bigwedge i. i \in I \implies (\lambda x. f \ (i,x)) \in \text{borel-measurable } (Mi \ i)$
and $(\sum n. (\int^+ x. \text{norm } (f \text{ (from-n } n, x)) \partial(\text{Mi (from-n } n)))) < \infty$
shows $\text{integrable (coPiM I Mi) f}$

using *nn-integral-coPiM-countable-infinite*[*OF assms(1), of - Mi*] *assms(2,3)*
by(*auto simp: measurable-copair-iff' integrable-iff-bounded*)

lemmas *integrable-coPiM-countable-infiniteI' = integrable-coPiM-countable-infiniteI*[*OF bij-betw-from-nat-into*]

lemmas *integrable-coPiM-natI = integrable-coPiM-countable-infiniteI*[*OF bij-id, simplified id-apply*]

lemma *integral-coPiM-countable-infinite*:

fixes $f :: - \Rightarrow 'b::\{\text{banach, second-countable-topology}\}$

assumes *bij-betw from-n (UNIV :: nat set) I integrable (coPiM I Mi) f*

shows $(\int x. f x \partial(\text{coPiM } I \text{ } Mi)) = (\sum n. (\int x. f (\text{from-n } n, x) \partial(Mi (\text{from-n } n))))$ (**is** *?lhs = ?rhs*)

proof –

have *?lhs* = $(\sum_{\infty} i \in I. (\int x. f (i, x) \partial Mi i))$

by(*simp add: integral-coPiM assms*)

also have ... = $(\sum_{\infty} i \in \text{from-n } ' UNIV. (\int x. f (i, x) \partial Mi i))$

by(*rule arg-cong[where f=infsum -] (metis assms(1) bij-betw-def)*)

also have ... = $(\sum_{\infty} n \in UNIV. (\int x. f (\text{from-n } n, x) \partial(Mi (\text{from-n } n))))$

by(*rule infsum-reindex[simplified comp-def] (use assms(1) bij-betw-imp-inj-on*

in blast)

also have ... = *?rhs*

by(*auto intro!: infsum-eq-suminf norm-summable-imp-summable-on integrable-coPiM-countable-infinite-sum assms*)

finally show *?thesis* .

qed

lemmas *integral-coPiM-countable-infinite' = integral-coPiM-countable-infinite*[*OF bij-betw-from-nat-into*]

lemmas *integral-coPiM-nat = integral-coPiM-countable-infinite*[*OF bij-id, simplified id-apply*]

3.8 Finiteness

lemma *finite-measure-coPiM*:

assumes *finite I $\wedge i. i \in I \implies$ finite-measure (Mi i)*

shows *finite-measure (coPiM I Mi)*

by(*rule finite-measureI (auto simp: emeasure-coPiM-finite finite-measure.emeasure-finite assms)*)

3.9 σ -Finiteness

lemma *sigma-finite-measure-coPiM*:

assumes *countable I $\wedge i. i \in I \implies$ sigma-finite-measure (Mi i)*

shows *sigma-finite-measure (coPiM I Mi)*

proof

have $\exists A. \text{range } A \subseteq \text{sets } (Mi i) \wedge (\bigcup n. A n) = \text{space } (Mi i) \wedge (\forall n::\text{nat}. \text{emeasure } (Mi i) (A n) \neq \infty)$

if $i \in I$ **for** i

using *sigma-finite-measure.sigma-finite*[*OF assms(2)[OF that]*] **by** *metis*

hence $\exists A. \forall i \in I. \text{range } (A \ i) \subseteq \text{sets } (M \ i) \wedge (\bigcup n. A \ i \ n) = \text{space } (M \ i) \wedge$
 $(\forall n::\text{nat}. \text{emeasure } (M \ i) (A \ i \ n) \neq \infty)$
by *metis*
then obtain A_i
where $A_i[\text{measurable}]: \bigwedge i \ n. i \in I \implies A_i \ i \ n \in \text{sets } (M \ i)$
 $\bigwedge i. i \in I \implies (\bigcup n::\text{nat}. (A_i \ i \ n)) = \text{space } (M \ i)$
 $\bigwedge i \ n. i \in I \implies \text{emeasure } (M \ i) (A_i \ i \ n) \neq \infty$
by (*metis UNIV-I sets-range*)
show $\exists A. \text{countable } A \wedge A \subseteq \text{sets } (\text{coPiM } I \ M) \wedge \bigcup A = \text{space } (\text{coPiM } I \ M)$
 $\wedge (\forall a \in A. \text{emeasure } (\text{coPiM } I \ M) a \neq \infty)$
proof (*intro exI[where x= $\bigcup n. (\bigcup i \in I. \{\{i\} \times A_i \ i \ n\})$] conjI ballI*)
show $\text{countable } (\bigcup n. (\bigcup i \in I. \{\{i\} \times A_i \ i \ n\}))$
using *assms(1) by auto*
next
show $(\bigcup n. \bigcup i \in I. \{\{i\} \times A_i \ i \ n\}) \subseteq \text{sets } (\text{coPiM } I \ M)$
by *auto*
next
show $\bigcup (\bigcup n. \bigcup i \in I. \{\{i\} \times A_i \ i \ n\}) = \text{space } (\text{coPiM } I \ M)$
using *sets.sets-into-space[OF Ai(1)] Ai(2) by(fastforce simp: space-coPiM)*
next
fix a
assume $a \in (\bigcup n. \bigcup i \in I. \{\{i\} \times A_i \ i \ n\})$
then obtain $n \ i$ **where** $a: i \in I \ a = \{i\} \times A_i \ i \ n$
by *blast*
show $\text{emeasure } (\text{coPiM } I \ M) a \neq \infty$
using $a(1) \ Ai(3) \ \text{assms}$ **by** (*auto simp: a(2) emeasure-coPiM-coproj*)
qed
qed
end

4 Additional Properties

theory *Coproduct-Measure-Additional*
imports *Coproduct-Measure*
Standard-Borel-Spaces.StandardBorel
S-Finite-Measure-Monad.Kernels
S-Finite-Measure-Monad.Measure-QuasiBorel-Adjunction
begin

4.1 S-Finiteness

lemma *s-finite-measure-copair-measure*:
assumes *s-finite-measure* M *s-finite-measure* N
shows *s-finite-measure* (*copair-measure* $M \ N$)
proof –
note $[\text{measurable}] = \text{measurable-vimage-Inl}[of \ - \ M \ N] \ \text{measurable-vimage-Inr}[of \ - \ M \ N]$
obtain $M_i \ N_i$ **where** $[\text{measurable-cong}]$:

$\bigwedge i. \text{sets } (M_i i) = \text{sets } M \bigwedge i. \text{finite-measure } (M_i i) \bigwedge A. M A = (\sum i. M_i i A)$
 $\bigwedge i. \text{sets } (N_i i) = \text{sets } N \bigwedge i. \text{finite-measure } (N_i i) \bigwedge A. N A = (\sum i. N_i i A)$
by (*metis* *assms*(1) *assms*(2) *s-finite-measure.finite-measures'*)
thus *?thesis*
by(*auto intro!*: *s-finite-measureI*[**where** $M_i = \lambda i. M_i i \oplus_M N_i i$] *finite-measure-copair-measure*
cong: sets-copair-measure-cong simp: emeasure-copair-measure sum-
inf-add)
qed

lemma *s-finite-measure-coPiM*:
assumes *countable I* $\bigwedge i. i \in I \implies \text{s-finite-measure } (M_i i)$
shows *s-finite-measure (coPiM I Mi)*
proof –
note *measurable-Pair-vimage*[*measurable (raw)*]
consider *finite I | infinite I countable I*
using *assms* **by** *argo*
then show *?thesis*
proof cases
assume *I:finite I*
show *?thesis*
by(*auto intro!*: *s-finite-measure-finite-sumI*[**where** $M_i = \lambda i. \text{distr } (M_i i) (\text{coPiM } I \text{ Mi}) (\text{Pair } i)$
 $\text{and } I = I, OF - \text{s-finite-measure.s-finite-measure-distr}[OF$
 $\text{assms}(2)]$
simp: emeasure-distr emeasure-coPiM-finite I)
next
assume *I:infinite I countable I*
then have [*simp*]: $\bigwedge n. \text{from-nat-into } I \ n \in I$
by (*simp add: from-nat-into infinite-imp-nonempty*)
show *?thesis*
by(*auto intro!*: *s-finite-measure-s-finite-sumI*[**where**
 $M_i = \lambda n. \text{distr } (M_i (\text{from-nat-into } I \ n)) (\text{coPiM } I \text{ Mi}) (\text{Pair } (\text{from-nat-into } I \ n)),$
 $OF - \text{s-finite-measure.s-finite-measure-distr}[OF \ \text{assms}(2)]$
simp: emeasure-distr I emeasure-coPiM-countable-infinite' coPair-inverse-space-unit[**where**
 $I = I$])
qed
qed

4.2 Standardness

lemma *standard-borel-copair-measure*:
assumes *standard-borel M standard-borel N*
shows *standard-borel (M \oplus_M N)*
proof –
obtain *A* **where** $A[\text{measurable}]: A \in \text{sets borel } A \subseteq \{0 < .. < 1 :: \text{real}\}$
 $M \text{ measurable-isomorphic restrict-space borel } A$
by (*meson* *assms*(1) *greaterThanLessThan-borel linorder-not-le not-one-le-zero*
standard-borel.isomorphic-subset-real uncountable-open-interval)

```

then obtain  $f f'$ 
  where  $f$ [measurable]:  $f \in M \rightarrow_M \text{restrict-space borel } A$ 
     $f' \in \text{restrict-space borel } A \rightarrow_M M$ 
     $\bigwedge x. x \in \text{space } M \implies f' (f x) = x \wedge y. y \in A \implies f (f' y) = y$ 
  using measurable-isomorphicD[OF A(3)] unfolding space-restrict-space by
fastforce
obtain  $B$  where  $B$ [measurable]:  $B \in \text{sets borel } B \subseteq \{1 <..<2::\text{real}\}$ 
     $N \text{ measurable-isomorphic restrict-space borel } B$ 
by (metis assms(2) greaterThanLessThan-borel linorder-not-le numeral-le-one-iff
semiring-norm(69) standard-borel.isomorphic-subset-real uncount-
able-open-interval)
then obtain  $g g'$ 
  where  $g$ [measurable]:  $g \in N \rightarrow_M \text{restrict-space borel } B$ 
     $g' \in \text{restrict-space borel } B \rightarrow_M N$ 
     $\bigwedge x. x \in \text{space } N \implies g' (g x) = x$ 
     $\bigwedge y. y \in B \implies g (g' y) = y$ 
  using measurable-isomorphicD[OF B(3)] unfolding space-restrict-space by
fastforce
have  $AB: A \cap B = \{\}$ 
  using  $A B$  by fastforce
have [measurable]:  $f \in M \rightarrow_M \text{restrict-space borel } (A \cup B)$ 
  using  $f(1)$  unfolding measurable-restrict-space2-iff by blast
have [measurable]:  $g \in N \rightarrow_M \text{restrict-space borel } (A \cup B)$ 
  using  $g(1)$  unfolding measurable-restrict-space2-iff by blast

have iso:  $\text{restrict-space borel } (A \cup B) \text{ measurable-isomorphic } M \oplus_M N$ 
proof (safe intro!: measurable-isomorphic-byWitness)
  show case-sum  $f g \in M \oplus_M N \rightarrow_M \text{restrict-space borel } (A \cup B)$ 
    by (auto intro!: measurable-copair-Inl-Inr)
  show ( $\lambda r. \text{if } r \in A \text{ then } \text{Inl } (f' r) \text{ else if } r \in B \text{ then } \text{Inr } (g' r) \text{ else undefined}$ )
     $\in \text{restrict-space borel } (A \cup B) \rightarrow_M M \oplus_M N$  (is  $?f \in -$ )
proof -
  have 1:
     $\text{restrict-space } (\text{restrict-space borel } (A \cup B)) \{r. r \in A\} = \text{restrict-space}$ 
    borel } A
     $\text{restrict-space } (\text{restrict-space borel } (A \cup B)) \{r. r \notin A\} = \text{restrict-space}$ 
    borel } B
     $\text{restrict-space } (\text{restrict-space borel } B) \{x. x \in B\} = \text{restrict-space borel } B$ 
     $\text{restrict-space } (\text{restrict-space borel } B) \{x. x \notin B\} = \text{count-space } \{\}$ 
  using  $AB$  by (auto simp: restrict-restrict-space
    intro!: arg-cong[where f=restrict-space borel] space-empty)
  have 2:  $\{r \in \text{space } (\text{restrict-space borel } (A \cup B)). r \in A\} = A$ 
     $\{x \in \text{space } (\text{restrict-space } (\text{restrict-space borel } (A \cup B)) \{r. r \notin A\}). x$ 
     $\in B\} = B$ 
     $\{x \in \text{space } (\text{restrict-space borel } B). x \in B\} = B$ 
  using  $AB$  by (auto simp: space-restrict-space)
show ?thesis
  by (intro measurable-If-restrict-space-iff[THEN iffD2] conjI)
    (unfold 1 2, simp-all add: sets-restrict-space-iff)

```

```

qed
show  $\bigwedge x. x \in \text{space } (M \oplus_M N) \implies ?f (\text{case-sum } f \ g \ x) = x$ 
       $\bigwedge r. r \in \text{space } (\text{restrict-space borel } (A \cup B)) \implies \text{case-sum } f \ g \ (?f \ r) = r$ 
      using measurable-space[OF f(1)] measurable-space[OF g(1)] AB
      by (auto simp: space-copair-measure f g)
qed
show ?thesis
      by(auto intro!: standard-borel.measurable-isomorphic-standard[OF - iso]
          standard-borel.standard-borel-restrict-space[OF standard-borel-ne.standard-borel])
qed

corollary
shows standard-borel-ne-copair-measure1: standard-borel-ne M  $\implies$  standard-borel
N  $\implies$  standard-borel-ne (M  $\oplus_M$  N)
  and standard-borel-ne-copair-measure2: standard-borel M  $\implies$  standard-borel-ne
N  $\implies$  standard-borel-ne (M  $\oplus_M$  N)
  and standard-borel-ne-copair-measure: standard-borel-ne M  $\implies$  standard-borel-ne
N  $\implies$  standard-borel-ne (M  $\oplus_M$  N)
  by(auto simp: standard-borel-ne-def standard-borel-ne-axioms-def standard-borel-copair-measure
space-copair-measure)

lemma standard-borel-coPiM:
assumes countable I  $\bigwedge i. i \in I \implies$  standard-borel (Mi i)
shows standard-borel (coPiM I Mi)
proof -
let ?I = {i ∈ I. space (Mi i) ≠ {}}
have countable-I: countable ?I
  using assms by auto
define I' where I'  $\equiv$  to-nat-on ?I ' ?I
define Mn where Mn  $\equiv$   $\lambda n. Mi$  (from-nat-into ?I n)
have I':countable I'  $\bigwedge n. n \in I' \implies$  space (Mn n)  $\neq$  {}
       $\bigwedge n. n \in I' \implies$  standard-borel-ne (Mn n)
  using countable-I from-nat-into-to-nat-on[OF countable-I] assms(2)
by(fastforce simp: I'-def Mn-def standard-borel-ne-def standard-borel-ne-axioms-def
simp del: from-nat-into-to-nat-on)+
have iso1:coPiM I Mi measurable-isomorphic coPiM I' Mn
proof(safe intro!: measurable-isomorphic-byWitness[where f= $\lambda(i,x).$  (to-nat-on
?I i, x)
      and g= $\lambda(n,x).$  (from-nat-into ?I n, x)])
show ( $\lambda(i, x).$  (to-nat-on ?I i, x))  $\in$  coPiM I Mi  $\rightarrow_M$  coPiM I' Mn
proof(rule measurable-coPiM2)
  fix i
  assume i:i ∈ I
  show Pair (to-nat-on ?I i)  $\in$  Mi i  $\rightarrow_M$  coPiM I' Mn
proof(cases space (Mi i) = {})
  assume space (Mi i)  $\neq$  {}
  then show ?thesis
    by(intro measurable-compose[OF - measurable-Pair-coPiM[where I=I']]
      (use I'-def i countable-I Mn-def in auto)

```

```

    qed(simp add: measurable-def)
  qed
  show  $(\lambda(n,x). (from-nat-into ?I n, x)) \in coPiM I' Mn \rightarrow_M coPiM I Mi$ 
  proof(rule measurable-coPiM2)
    fix n
    assume  $n \in I'$ 
    show Pair (from-nat-into ?I n)  $\in Mn n \rightarrow_M coPiM I Mi$ 
      by (metis (no-types, lifting) Mn-def I'-def  $\langle n \in I' \rangle$  emptyE empty-is-image
          from-nat-into measurable-Pair-coPiM mem-Collect-eq)
    qed
  qed(auto intro!: from-nat-into-to-nat-on to-nat-on-from-nat-into simp: space-coPiM
  I'-def countable-I)
  have  $\exists A. A \in sets \text{ borel} \wedge A \subseteq \{real\ n <.. real\ n + 1\} \wedge Mn\ n \text{ measurable-isomorphic (restrict-space borel } A)$ 
    if  $n:n \in I'$  for n
    using standard-borel.isomorphic-subset-real[OF
      standard-borel-ne.standard-borel[OF I'(3)[OF n]], of  $\{real\ n <.. real\ n + 1\}$ ]
      uncountable-half-open-interval-2[of real n real n + 1]
    by fastforce
  then obtain An'
    where An':  $\bigwedge n. n \in I' \implies An' n \in sets \text{ borel}$ 
       $\bigwedge n. n \in I' \implies An' n \subseteq \{real\ n <.. real\ n + 1\}$ 
       $\bigwedge n. n \in I' \implies Mn\ n \text{ measurable-isomorphic (restrict-space borel (An' n))}$ 
    by metis
  define An where  $An \equiv \lambda n. \text{if } n \in I' \text{ then } An' n \text{ else } \{real\ n + 1\}$ 
  have An[measurable]:  $\bigwedge n. An\ n \in sets \text{ borel}$ 
     $\bigwedge n. An\ n \subseteq \{real\ n <.. real\ n + 1\}$ 
     $\bigwedge n. n \in I' \implies Mn\ n \text{ measurable-isomorphic (restrict-space borel (An n))}$ 
  (An n))
  using An' by(auto simp: An-def)
  hence disj-An: disjoint-family An
  unfolding disjoint-family-on-def
  by safe (metis (no-types, opaque-lifting) greaterThanAtMost-iff less-le nat-less-real-le
  not-less order-trans subset-eq)
  obtain fn gn'
    where fg:  $\bigwedge n. n \in I' \implies fn\ n \in Mn\ n \rightarrow_M \text{restrict-space borel (An n)}$ 
       $\bigwedge n. n \in I' \implies gn' n \in \text{restrict-space borel (An n)} \rightarrow_M Mn\ n$ 
       $\bigwedge n\ x. n \in I' \implies x \in \text{space (Mn n)} \implies gn' n (fn n x) = x$ 
       $\bigwedge n\ r. n \in I' \implies r \in \text{space (restrict-space borel (An n))} \implies fn n (gn' n r) = r$ 
    using measurable-isomorphicD[OF An(3)] by metis
  define gn where  $gn \equiv (\lambda n\ r. \text{if } r \in An\ n \text{ then } gn' n r \text{ else (SOME } x. x \in \text{space (Mn n))})$ 
  have gn-meas[measurable]:  $gn\ n \in \text{borel} \rightarrow_M Mn\ n$  if  $n:n \in I'$  for n
  unfolding gn-def by(rule measurable-restrict-space-iff[THEN iffD1, OF - -
  fg(2)[OF n]])
  (auto simp add: I'(2) some-in-eq that)
  have fg':  $\bigwedge n\ x. n \in I' \implies x \in \text{space (Mn n)} \implies gn\ n (fn n x) = x$ 

```

```

       $\bigwedge n r. n \in I' \implies r \in An\ n \implies fn\ n\ (gn\ n\ r) = r$ 
    using fg measurable-space[OF fg(1)] by(auto simp: gn-def)
    have fn[measurable]:fn n ∈ Mn n →M restrict-space borel (⋃ n∈I'. An n) if n:n
    ∈ I' for n
      using measurable-restrict-space2-iff[THEN iffD1,OF fg(1)[OF n]]
      by(auto intro!: measurable-restrict-space2 n)
    let ?f = λ(n,x). fn n x and ?g = λr. (nat ⌈r⌉ - 1, gn (nat ⌈r⌉ - 1) r)
    have iso2:coPiM I' Mn measurable-isomorphic restrict-space borel (⋃ n∈I'. An
    n)
    proof(safe intro!: measurable-isomorphic-byWitness)
      show ?f ∈ coPiM I' Mn →M restrict-space borel (⋃ n∈I'. An n)
        by(auto intro!: measurable-coPiM2)
    next
      show ?g ∈ restrict-space borel (⋃ n∈I'. An n) →M coPiM I' Mn
    proof(safe intro!: measurable-coPiM1)
      have 1:restrict-space borel (⋃ (An ' I')) →M count-space I'
        = restrict-space borel (⋃ (An ' I')) →M restrict-space (count-space
    UNIV) I'
        by (simp add: restrict-count-space)
      show (λx. nat ⌈x⌉ - 1) ∈ restrict-space borel (⋃ (An ' I')) →M count-space
    I'
        unfolding 1
    proof(safe intro!: measurable-restrict-space3)
      fix n r
      assume n:n ∈ I' r ∈ An n
      then have real n < r r ≤ real n + 1
        using An(2) by fastforce+
      thus nat ⌈r⌉ - 1 ∈ I'
        by (metis n(1) add.commute diff-Suc-1 le-SucE nat-ceiling-le-eq not-less
    of-nat-Suc)
      qed simp
    qed(auto simp: measurable-restrict-space1)
    next
      fix n x
      assume (n,x)∈space (coPiM I' Mn)
      then have nx:n ∈ I' x ∈ space (Mn n)
        by(auto simp: space-coPiM)
      have 1:nat ⌈?f (n,x)⌉ = n + 1
        using measurable-space[OF fg(1)[OF nx(1)] nx(2)] An(2)[of n]
        by simp
        (metis add.commute greaterThanAtMost-iff le-SucE nat-ceiling-le-eq not-less
    of-nat-Suc subset-eq)
      show ?g (?f (n,x)) = (n,x)
        unfolding 1 using fg'(1)[OF nx] by simp
    next
      fix y
      assume y ∈ space (restrict-space borel (⋃ (An ' I)))
      then obtain n where n: n ∈ I' y ∈ An n
        by auto

```



```

then have [simp]: nat [y] = n + 1
using An(2)[of n]
by simp (metis add commute greaterThanAtMost-iff le-SucE nat-ceiling-le-eq
not-less-of-nat-Suc subset-eq)
show ?f (?g y) = y
using fg'(2)[OF n(1)] n(2) by auto
qed
have standard-borel (restrict-space borel (⋃ (An ' I)))
by(auto intro!: standard-borel-ne.standard-borel[THEN standard-borel.standard-borel-restrict-space])
with iso1 iso2 show ?thesis
by (meson measurable-isomorphic-sym standard-borel.measurable-isomorphic-standard)
qed

```

```

lemma standard-borel-ne-coPiM:
assumes countable I  $\wedge i. i \in I \implies$  standard-borel (Mi i)
and  $i \in I$  space (Mi i)  $\neq \{\}$ 
shows standard-borel-ne (coPiM I Mi)
proof -
have space (coPiM I Mi)  $\neq \{\}$ 
using assms(3) assms(4) space-coPiM by fastforce
thus ?thesis
by(auto intro!: standard-borel-coPiM assms simp: standard-borel-ne-def stan-
dard-borel-ne-axioms-def)
qed

```

4.3 Relationships with Quasi-Borel Spaces

Proposition19(3) [1]

```

lemma r-preserve-copair: measure-to-qbs (copair-measure M N) = measure-to-qbs
M  $\oplus_Q$  measure-to-qbs N
proof(safe intro!: qbs-eqI)
fix  $\alpha$ 
assume  $\alpha \in$  qbs-Mx (measure-to-qbs (M  $\oplus_M$  N))
then have a[measurable]:  $\alpha \in$  borel  $\rightarrow_M$  M  $\oplus_M$  N
by(simp add: qbs-Mx-R)
have s[measurable]:  $\alpha - ' Inr ' space N \in$  sets borel  $\alpha - ' Inl ' space M \in$  sets
borel
by(auto intro!: measurable-sets-borel[OF a])
consider  $\alpha - ' Inl ' space M \cap space borel = space borel$ 
|  $\alpha - ' Inr ' (space N) \cap space borel = space borel$ 
|  $\alpha - ' Inl ' space M \cap space borel \subset space borel$ 
|  $\alpha - ' Inr ' (space N) \cap space borel \subset space borel$ 
by blast
then show  $\alpha \in$  qbs-Mx (measure-to-qbs M  $\oplus_Q$  measure-to-qbs N)
proof cases
assume 1:  $\alpha - ' Inl ' space M \cap space borel = space borel$ 
then obtain f' where f'  $\in$  borel  $\rightarrow_M$  M  $\wedge x. x \in space borel \implies \alpha x = Inl$ 
(f' x)
using measurable-copair-dest1[OF a] by blast

```

```

thus ?thesis
  using 1 by(auto simp: copair-qbs-Mx copair-qbs-Mx-def qbs-Mx-R
    intro!: bexI[where x=α - 'Inr ' space N] bexI[where x=f'])
next
  assume 2:α - 'Inr ' space N ∩ space borel = space borel
  then obtain f' where f' ∈ borel →M N ∧ x. x ∈ space borel ⇒ α x = Inr
(f' x)
  using measurable-copair-dest2[OF a] by blast
  thus ?thesis
  using 2 by(auto simp: copair-qbs-Mx copair-qbs-Mx-def qbs-Mx-R
    intro!: bexI[where x=α - 'Inr ' space N] bexI[where x=f'])
next
  case 3
  then obtain f' f''
  where f[measurable]:f' ∈ borel →M M
    f'' ∈ borel →M N
    ∧ x. x ∈ space borel ⇒ x ∈ α - 'Inl ' space M ⇒ α x =
Inl (f' x)
    ∧ x. x ∈ space borel ⇒ x ∉ α - 'Inl ' space M ⇒ α x =
Inr (f'' x)
  using measurable-copair-dest3[OF a] by metis
  moreover have α - 'Inl ' space M ≠ UNIV α - 'Inl ' space M ≠ {}
  using 3 measurable-space[OF a] by(fastforce simp: space-copair-measure)+
  ultimately show ?thesis
  by(auto simp: copair-qbs-Mx copair-qbs-Mx-def qbs-Mx-R simp del: vimage-eq
    intro!: bexI[where x=α - 'Inl ' space M] bexI[where x=f'] bexI[where
x=f''])
  qed
qed(auto simp: qbs-Mx-R copair-qbs-Mx copair-qbs-Mx-def)

lemma r-preserve-coproduct:
  assumes countable I
  shows measure-to-qbs (coPiM I M) = (∏Q i∈I. measure-to-qbs (M i))
proof(safe intro!: qbs-eqI)
  fix α
  assume h:α ∈ qbs-Mx (measure-to-qbs (coPiM I M))
  then obtain a g
  where a ∈ borel →M count-space I
    ∧ i. i ∈ I ⇒ space (M i) ≠ {} ⇒ g i ∈ borel →M M i
    α = (λx. (a x, g (a x) x))
  using measurable-coPiM1-elements[OF assms] unfolding qbs-Mx-R by blast
  thus α ∈ qbs-Mx (∏Q i∈I. measure-to-qbs (M i))
  using qbs-Mx-to-X[OF h]
  by(safe intro!: coPiQ-MxI) (auto simp: qbs-Mx-R qbs-space-R space-coPiM)
next
  fix α
  assume α ∈ qbs-Mx (∏Q i∈I. measure-to-qbs (M i))
  then obtain a g where a ∈ borel →M count-space I
    ∧ i. i ∈ range a ⇒ g i ∈ borel →M M i α = (λx. (a x, g (a

```

```

x) x))
  unfolding coPiQ-Mx coPiQ-Mx-def qbs-Mx-R by blast
  thus  $\alpha \in$  qbs-Mx (measure-to-qbs (coPiM I M))
  by(auto intro!: measurable-coPiM1' simp: qbs-Mx-R assms)
qed

end

```

References

- [1] C. Heunen, O. Kammar, S. Staton, and H. Yang. A convenient category for higher-order probability theory. In *Proceedings of the 32nd Annual ACM/IEEE Symposium on Logic in Computer Science, LICS '17*. IEEE Press, 2017.