

Concentration Inequalities

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Abstract

Concentration inequalities provide bounds on how a random variable (or a sum/composition of random variables) deviate from their expectation, usually based on moments/independence of the variables.

The most important concentration inequalities (the Markov, Chebyshev, and Hoelder inequalities and the Chernoff bounds) are already part of HOL-Probability. This entry collects more advanced results, such as Bennett's/Bernstein's Inequality, Bienaymé's Identity, Cantelli's Inequality, the Efron-Stein Inequality, McDiarmid's Inequality, and the Paley-Zygmund Inequality.

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1 Preliminary results

```
theory Concentration-Inequalities-Preliminary
  imports Lp.Lp
begin
```

Version of Cauchy-Schwartz for the Lebesgue integral:

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lemma *cauchy-schwartz*:
fixes $f\ g :: - \Rightarrow \text{real}$
assumes $f \in \text{borel-measurable } M\ g \in \text{borel-measurable } M$
assumes $\text{integrable } M (\lambda x. (f\ x) \wedge 2)\ \text{integrable } M (\lambda x. (g\ x) \wedge 2)$
shows $\text{integrable } M (\lambda x. f\ x * g\ x)$ (**is** ?A)
 $(\int x. f\ x * g\ x\ \partial M) \leq (\int x. (f\ x) \wedge 2\ \partial M) \text{ powr } (1/2) * (\int x. (g\ x) \wedge 2\ \partial M) \text{ powr } (1/2)$
(is ?L \leq ?R)
proof –
show 0: ?A
using *assms* **by** (*intro Holder-inequality(1)*[**where** $p=2$ **and** $q=2$])
auto

have ?L $\leq (\int x. |f\ x * g\ x|\ \partial M)$
using 0 **by** (*intro integral-mono*) *auto*
also have ... $\leq (\int x. |f\ x| \text{ powr } 2\ \partial M) \text{ powr } (1/2) * (\int x. |g\ x| \text{ powr } 2\ \partial M) \text{ powr } (1/2)$
using *assms* **by** (*intro Holder-inequality(2)*) *auto*
also have ... = ?R **by** *simp*
finally show ?L \leq ?R **by** *simp*
qed

Generalization of *prob-space.indep-vars-iff-distr-eq-PiM'*:

lemma (**in** *prob-space*) *indep-vars-iff-distr-eq-PiM''*:
fixes $I :: 'i \text{ set}$ **and** $X :: 'i \Rightarrow 'a \Rightarrow 'b$
assumes $rv: \bigwedge i. i \in I \implies \text{random-variable } (M' i) (X i)$
shows $\text{indep-vars } M' X I \longleftrightarrow$
 $\text{distr } M (\prod_M i \in I. M' i) (\lambda x. \lambda i \in I. X i x) = (\prod_M i \in I. \text{distr } M (M' i) (X i))$
proof (*cases* $I = \{\}$)
case *True*
have 0: $\text{indicator } A (\lambda -. \text{undefined}) = \text{emeasure } (\text{count-space } \{\lambda -. \text{undefined}\}) A$ (**is** ?L = ?R)
if $A \subseteq \{\lambda -. \text{undefined}\}$ **for** $A :: ('i \Rightarrow 'b) \text{ set}$
proof –
have 1: $A \neq \{\} \implies A = \{\lambda -. \text{undefined}\}$
using *that* **by** *auto*

have ?R = *of-nat* (*card* A)
using *finite-subset that* **by** (*intro emeasure-count-space-finite that*)
auto
also have ... = ?L
using 1 **by** (*cases* $A = \{\}$) *auto*
finally show ?thesis **by** *simp*
qed

have $\text{distr } M (\prod_M i \in I. M' i) (\lambda x. \lambda i \in I. X i x) =$
 $\text{distr } M (\text{count-space } \{\lambda -. \text{undefined}\}) (\lambda -. (\lambda -. \text{undefined}))$
unfolding *True PiM-empty* **by** (*intro distr-cong*) (*auto simp: restrict-def*)

also have ... = return (count-space {λ-. undefined}) (λ-. undefined)
by (intro distr-const) auto
also have ... = count-space ({λ-. undefined} :: ('i ⇒ 'b) set)
by (intro measure-eqI) (auto simp:0)
also have ... = (Π_M i∈I. distr M (M' i) (X i))
unfolding True PiM-empty **by** simp
finally have distr M (Π_M i∈I. M' i) (λx. λi∈I. X i x) = (Π_M i∈I.
distr M (M' i) (X i)) ↔ True
by simp
also have ... ↔ indep-vars M' X I
unfolding indep-vars-def **by** (auto simp add: space-PiM indep-sets-def)
(auto simp add: True)
finally show ?thesis **by** simp
next
case False
thus ?thesis
by (intro indep-vars-iff-distr-eq-PiM' assms) auto
qed

lemma proj-indep:

assumes $\bigwedge i. i \in I \implies \text{prob-space } (M i)$
shows prob-space.indep-vars (PiM I M) M (λi ω. ω i) I

proof –

interpret prob-space (PiM I M)
by (intro prob-space-PiM assms)

have distr (Pi_M I M) (Pi_M I M) (λx. restrict x I) = PiM I M
by (intro distr-PiM-reindex assms) auto

also have ... = Pi_M I (λi. distr (Pi_M I M) (M i) (λω. ω i))
by (intro PiM-cong refl distr-PiM-component[symmetric] assms)

finally have
distr (Pi_M I M) (Pi_M I M) (λx. restrict x I) = Pi_M I (λi. distr
(Pi_M I M) (M i) (λω. ω i))
by simp

thus indep-vars M (λi ω. ω i) I
by (intro iffD2[OF indep-vars-iff-distr-eq-PiM'']) simp-all

qed

lemma forall-Pi-to-PiE:

assumes $\bigwedge x. P x = P (\text{restrict } x I)$
shows $(\forall x \in \text{Pi } I A. P x) = (\forall x \in \text{PiE } I A. P x)$
using assms **by** (simp add: PiE-def Pi-def set-eq-iff, force)

lemma PiE-reindex:

assumes inj-on f I
shows PiE I (A ∘ f) = (λa. restrict (a ∘ f) I) ‘ PiE (f ‘ I) A (is
?lhs = ?g ‘ ?rhs)

proof –

have ?lhs ⊆ ?g ‘ ?rhs

proof (*rule subsetI*)
fix x
assume $a: x \in \text{PiE } I (A \circ f)$
define y **where** $y\text{-def}: y = (\lambda k. \text{if } k \in f \text{ ' } I \text{ then } x \text{ (the-inv-into } I$
 $f k) \text{ else undefined})$
have $b: y \in \text{PiE } (f \text{ ' } I) A$
using a *assms the-inv-into-f-eq[OF assms]*
by (*simp add: y-def PiE-iff extensional-def*)
have $c: x = (\lambda a. \text{restrict } (a \circ f) I) y$
using a *assms the-inv-into-f-eq extensional-arb*
by (*intro ext, simp add: y-def PiE-iff, fastforce*)
show $x \in ?g \text{ ' } ?rhs$ **using** $b c$ **by** *blast*
qed
moreover **have** $?g \text{ ' } ?rhs \subseteq ?lhs$
by (*rule image-subsetI, simp add: Pi-def PiE-def*)
ultimately **show** $?thesis$ **by** *blast*
qed

context *prob-space*
begin

lemma *indep-sets-reindex:*

assumes *inj-on f I*
shows *indep-sets A (f ' I) = indep-sets ($\lambda i. A (f i)$) I*
proof –
have $a: \bigwedge J. J \subseteq I \implies (\prod j \in f \text{ ' } J. g j) = (\prod j \in J. g (f j))$
by (*metis assms prod.reindex-cong subset-inj-on*)

have $b: J \subseteq I \implies (\prod_E i \in J. A (f i)) = (\lambda a. \text{restrict } (a \circ f) J) \text{ ' } A$
 $\text{PiE } (f \text{ ' } J) A$ **for** J
using *assms inj-on-subset*
by (*subst PiE-reindex[symmetric] auto*)

have $c: \bigwedge J. J \subseteq I \implies \text{finite } (f \text{ ' } J) = \text{finite } J$
by (*meson assms finite-image-iff inj-on-subset*)

show $?thesis$
by (*simp add: indep-sets-def all-subset-image a c*) (*simp-all add: forall-Pi-to-PiE*
 b)
qed

lemma *indep-vars-reindex:*

assumes *inj-on f I*
assumes *indep-vars M' X' (f ' I)*
shows *indep-vars (M' o f) ($\lambda k \omega. X' (f k) \omega$) I*
using *assms* **by** (*simp add: indep-vars-def2 indep-sets-reindex*)

lemma *indep-vars-cong-AE:*

assumes *AE x in M. ($\forall i \in I. X' i x = Y' i x$)*

assumes *indep-vars* $M' X' I$
assumes $\bigwedge i. i \in I \implies \text{random-variable } (M' i) (Y' i)$
shows *indep-vars* $M' Y' I$
proof –
have $a: AE\ x\ \text{in}\ M. (\lambda i \in I. Y' i\ x) = (\lambda i \in I. X' i\ x)$
by (*rule* $AE\text{-mp}[OF\ \text{assms}(1)]$, *rule* $AE\text{-I2}$, *simp cong:restrict-cong*)
have $b: \bigwedge i. i \in I \implies \text{random-variable } (M' i) (X' i)$
using $\text{assms}(2)$ **by** (*simp add:indep-vars-def2*)
have $c: \bigwedge x. x \in I \implies AE\ xa\ \text{in}\ M. X' x\ xa = Y' x\ xa$
by (*rule* $AE\text{-mp}[OF\ \text{assms}(1)]$, *rule* $AE\text{-I2}$, *simp*)

have $\text{distr } M (Pi_M\ I\ M') (\lambda x. \lambda i \in I. Y' i\ x) = \text{distr } M (Pi_M\ I\ M') (\lambda x. \lambda i \in I. X' i\ x)$
by (*intro distr-cong-AE measurable-restrict a b assms(3)*) *auto*
also have $\dots = Pi_M\ I (\lambda i. \text{distr } M (M' i) (X' i))$
using $\text{assms } b$ **by** (*subst indep-vars-iff-distr-eq-PiM''[symmetric]*)
auto
also have $\dots = Pi_M\ I (\lambda i. \text{distr } M (M' i) (Y' i))$
by (*intro PiM-cong distr-cong-AE c assms(3) b*) *auto*
finally have $\text{distr } M (Pi_M\ I\ M') (\lambda x. \lambda i \in I. Y' i\ x) = Pi_M\ I (\lambda i. \text{distr } M (M' i) (Y' i))$
by *simp*

thus *?thesis*
using $\text{assms}(3)$
by (*subst indep-vars-iff-distr-eq-PiM''*) *auto*
qed

end

Integrability of bounded functions on finite measure spaces:

lemma *bounded-const*: $\text{bounded } ((\lambda x. (c::\text{real})) ' T)$
by (*intro boundedI[where B=norm c]*) *auto*

lemma *bounded-exp*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes $\text{bounded } ((\lambda x. f\ x) ' T)$
shows $\text{bounded } ((\lambda x. \exp (f\ x)) ' T)$

proof –
obtain m **where** $\text{norm } (f\ x) \leq m$ **if** $x \in T$ **for** x
using assms **unfolding** *bounded-iff* **by** *auto*

thus *?thesis*
by (*intro boundedI[where B=exp m]*) *fastforce*
qed

lemma *bounded-mult-comp*:
fixes $f :: 'a \Rightarrow \text{real}$
assumes $\text{bounded } (f ' T)$ $\text{bounded } (g ' T)$

shows *bounded* $((\lambda x. (f x) * (g x)) \text{ ' } T)$
proof –
obtain $m1$ **where** $norm (f x) \leq m1$ $m1 \geq 0$ **if** $x \in T$ **for** x
using *assms unfolding bounded-iff* **by** *fastforce*
moreover obtain $m2$ **where** $norm (g x) \leq m2$ $m2 \geq 0$ **if** $x \in T$
for x
using *assms unfolding bounded-iff* **by** *fastforce*

ultimately show *?thesis*
by (*intro boundedI*[**where** $B=m1 * m2$]) (*auto intro!*: *mult-mono simp:abs-mult*)
qed

lemma *bounded-sum*:
fixes $f :: 'i \Rightarrow 'a \Rightarrow real$
assumes *finite I*
assumes $\bigwedge i. i \in I \implies bounded (f i \text{ ' } T)$
shows *bounded* $((\lambda x. (\sum i \in I. f i x)) \text{ ' } T)$
using *assms* **by** (*induction I*) (*auto intro:bounded-plus-comp bounded-const*)

lemma *bounded-pow*:
fixes $f :: 'a \Rightarrow real$
assumes *bounded* $((\lambda x. f x) \text{ ' } T)$
shows *bounded* $((\lambda x. (f x)^{\hat{n}}) \text{ ' } T)$
proof –
obtain m **where** $norm (f x) \leq m$ **if** $x \in T$ **for** x
using *assms unfolding bounded-iff* **by** *auto*
hence $norm ((f x)^{\hat{n}}) \leq m^{\hat{n}}$ **if** $x \in T$ **for** x
using *that unfolding norm-power* **by** (*intro power-mono*) *auto*
thus *?thesis* **by** (*intro boundedI*[**where** $B=m^{\hat{n}}$]) *auto*
qed

lemma *bounded-sum-list*:
fixes $f :: 'i \Rightarrow 'a \Rightarrow real$
assumes $\bigwedge y. y \in set\ ys \implies bounded (f y \text{ ' } T)$
shows *bounded* $((\lambda x. (\sum y \leftarrow ys. f y x)) \text{ ' } T)$
using *assms* **by** (*induction ys*) (*auto intro:bounded-plus-comp bounded-const*)

lemma (*in finite-measure*) *bounded-int*:
fixes $f :: 'i \Rightarrow 'a \Rightarrow real$
assumes *bounded* $((\lambda x. f (fst x) (snd x)) \text{ ' } (T \times space\ M))$
shows *bounded* $((\lambda x. (\int \omega. (f x \omega) \partial M)) \text{ ' } T)$
proof –
obtain m **where** $\bigwedge x y. x \in T \implies y \in space\ M \implies norm (f x y) \leq m$
using *assms unfolding bounded-iff* **by** *auto*
hence $m: \bigwedge x y. x \in T \implies y \in space\ M \implies norm (f x y) \leq max\ m\ 0$
by *fastforce*

```

have norm (∫ ω. (f x ω) ∂M) ≤ max m 0 * measure M (space M)
(is ?L ≤ ?R) if x ∈ T for x
proof –
  have ?L ≤ (∫ ω. norm (f x ω) ∂M) by simp
  also have ... ≤ (∫ ω. max m 0 ∂M)
    using that m by (intro integral-mono') auto
  also have ... = ?R
    by simp
  finally show ?thesis by simp
qed
thus ?thesis
  by (intro boundedI[where B=max m 0 * measure M (space M)])
auto
qed

```

```

lemmas bounded-intros =
  bounded-minus-comp bounded-plus-comp bounded-mult-comp bounded-sum
  finite-measure.bounded-int
  bounded-const bounded-exp bounded-pow bounded-sum-list

```

```

lemma (in prob-space) integrable-bounded:
  fixes f :: - ⇒ ('b :: {banach,second-countable-topology})
  assumes bounded (f ' space M)
  assumes f ∈ M →M borel
  shows integrable M f
proof –
  obtain m where norm (f x) ≤ m if x ∈ space M for x
    using assms(1) unfolding bounded-iff by auto
  thus ?thesis
    by (intro integrable-const-bound[where B=m] AE-I2 assms(2))
qed

```

```

lemma integrable-bounded-pmf:
  fixes f :: - ⇒ ('b :: {banach,second-countable-topology})
  assumes bounded (f ' set-pmf M)
  shows integrable (measure-pmf M) f
proof –
  obtain m where norm (f x) ≤ m if x ∈ set-pmf M for x
    using assms(1) unfolding bounded-iff by auto
  thus ?thesis by (intro measure-pmf.integrable-const-bound[where
  B=m] AE-pmfI) auto
qed

```

end

2 Bennett's Inequality

In this section we verify Bennett's inequality [1] and a (weak) version of Bernstein's inequality as a corollary. Both inequalities give concentration bounds for sums of independent random variables. The statement and proofs follow a summary paper by Boucheron et al. [2].

theory *Bennett-Inequality*

imports *Concentration-Inequalities-Preliminary*

begin

context *prob-space*

begin

lemma *indep-vars-Chernoff-ineq-ge:*

assumes *I: finite I*

assumes *ind: indep-vars (λ -. borel) X I*

assumes *sge: $s \geq 0$*

assumes *int: $\bigwedge i. i \in I \implies$ integrable $M (\lambda x. \exp (s * X i x))$*

shows *prob $\{x \in$ space $M. (\sum i \in I. X i x - \text{expectation } (X i)) \geq t\} \leq$*

*$\exp (-s*t) *$*

*$(\prod i \in I. \text{expectation } (\lambda x. \exp(s * (X i x - \text{expectation } (X i))))$*

proof (*cases $s = 0$*)

case [*simp*]: *True*

thus *?thesis*

by (*simp add: prob-space*)

next

case *False*

then have *s: $s > 0$ using sge by auto*

have [*measurable*]: *$\bigwedge i. i \in I \implies$ random-variable borel $(X i)$*

using *ind unfolding indep-vars-def by blast*

have *indep1: indep-vars (λ -. borel)*

*$(\lambda i \omega. \exp (s * (X i \omega - \text{expectation } (X i)))) I$*

apply (*intro indep-vars-compose[OF ind, unfolded o-def]*)

by *auto*

define *S where $S = (\lambda x. (\sum i \in I. X i x - \text{expectation } (X i)))$*

have *int1: $\bigwedge i. i \in I \implies$*

*integrable $M (\lambda \omega. \exp (s * (X i \omega - \text{expectation } (X i))))$*

by (*auto simp add: algebra-simps exp-diff int*)

have *eprod: $\bigwedge x. \exp (s * S x) = (\prod i \in I. \exp(s * (X i x - \text{expectation } (X i))))$*

unfolding *S-def*
by (*simp add: assms(1) exp-sum vector-space-over-itself.scale-sum-right*)

from *indep-vars-integrable[OF I indep1 int1]*
have *intS: integrable M (λx. exp (s * S x))*
unfolding *eprod by auto*

then have *si: set-integrable M (space M) (λx. exp (s * S x))*
unfolding *set-integrable-def*
apply (*intro integrable-mult-indicator*)
by *auto*

from *Chernoff-ineq-ge[OF s si]*
have *prob {x ∈ space M. S x ≥ t} ≤ exp (- s * t) * (∫ x∈space M. exp (s * S x) ∂M)*
by *auto*

also have ($\int x \in \text{space } M. \exp (s * S x) \partial M = \text{expectation } (\lambda x. \exp (s * S x))$)
unfolding *set-integral-space[OF intS] by auto*

also have $\dots = \text{expectation } (\lambda x. \prod_{i \in I}. \exp (s * (X i x - \text{expectation } (X i))))$
unfolding *S-def*
by (*simp add: assms(1) exp-sum vector-space-over-itself.scale-sum-right*)
also have $\dots = (\prod_{i \in I}. \text{expectation } (\lambda x. \exp (s * (X i x - \text{expectation } (X i))))$
apply (*intro indep-vars-lebesgue-integral[OF I indep1 int1]*) .
finally show *?thesis*
unfolding *S-def*
by *auto*

qed

definition *bennett-h::real ⇒ real*
where *bennett-h u = (1 + u) * ln (1 + u) - u*

lemma *exp-sub-two-terms-eq:*
fixes *x :: real*
shows $\exp x - x - 1 = (\sum n. x^{\wedge}(n+2) / \text{fact } (n+2))$
 $\text{summable } (\lambda n. x^{\wedge}(n+2) / \text{fact } (n+2))$
proof –
have ($\sum_{i < 2}. \text{inverse } (\text{fact } i) * x^{\wedge} i = 1 + x$)
by (*simp add:numeral-eq-Suc*)
thus $\exp x - x - 1 = (\sum n. x^{\wedge}(n+2) / \text{fact } (n+2))$
unfolding *exp-def*
apply (*subst suminf-split-initial-segment[where k = 2]*)
by (*auto simp add: summable-exp divide-inverse-commute*)
have *summable (λn. x^ \wedge n / fact n)*
by (*simp add: divide-inverse-commute summable-exp*)

then have $\text{summable } (\lambda n. x^{\wedge}(\text{Suc } (\text{Suc } n)) / \text{fact } (\text{Suc } (\text{Suc } n)))$
apply (*subst summable-Suc-iff*)
apply (*subst summable-Suc-iff*)
by *auto*
thus $\text{summable } (\lambda n. x^{\wedge}(n+2) / \text{fact } (n+2))$ **by** *auto*
qed

lemma *psi-mono*:

defines $f \equiv (\lambda x. (\exp x - x - 1) - x^{\wedge}2 / 2)$

assumes $xy: a \leq (b::\text{real})$

shows $f a \leq f b$

proof –

have 1: (*f has-real-derivative* ($\exp x - x - 1$)) (*at x*) **for** x

unfolding *f-def*

by (*auto intro!*; *derivative-eq-intros*)

have 2: $\bigwedge x. x \in \{a..b\} \implies 0 \leq \exp x - x - 1$

by (*smt (verit) exp-ge-add-one-self*)

from *deriv-nonneg-imp-mono*[*OF 1 2 xy*]

show *?thesis* **by** *auto*

qed

lemma *psi-inequality*:

assumes $le: x \leq (y::\text{real}) y \geq 0$

shows $y^{\wedge}2 * (\exp x - x - 1) \leq x^{\wedge}2 * (\exp y - y - 1)$

proof –

have $x: \exp x - x - 1 = (\sum n. (x^{\wedge}(n+2) / \text{fact } (n+2)))$

summable ($\lambda n. x^{\wedge}(n+2) / \text{fact } (n+2)$)

using *exp-sub-two-terms-eq* .

have $y: \exp y - y - 1 = (\sum n. (y^{\wedge}(n+2) / \text{fact } (n+2)))$

summable ($\lambda n. y^{\wedge}(n+2) / \text{fact } (n+2)$)

using *exp-sub-two-terms-eq* .

have $l: y^{\wedge}2 * (\exp x - x - 1) = (\sum n. y^{\wedge}2 * (x^{\wedge}(n+2) / \text{fact } (n+2)))$

using x

apply (*subst suminf-mult*)

by *auto*

have $ls: \text{summable } (\lambda n. y^{\wedge}2 * (x^{\wedge}(n+2) / \text{fact } (n+2)))$

by (*intro summable-mult*[*OF x(2)*])

have $r: x^{\wedge}2 * (\exp y - y - 1) = (\sum n. x^{\wedge}2 * (y^{\wedge}(n+2) / \text{fact } (n+2)))$

using y

```

apply (subst suminf-mult)
by auto
have rs: summable ( $\lambda n. x^2 * (y^{(n+2)} / \text{fact } (n+2))$ )
by (intro summable-mult[OF y(2)])

have  $|x| \leq |y| \vee |y| < |x|$  by auto
moreover {
  assume  $|x| \leq |y|$ 
  then have  $x^n \leq y^n$  for  $n$ 
  by (smt (verit, ccfv-threshold) bot-nat-0.not-eq-extremum le power-0
real-root-less-mono real-root-power-cancel root-abs-power)
  then have  $(x^2 * y^2) * x^n \leq (x^2 * y^2) * y^n$  for  $n$ 
  by (simp add: mult-left-mono)
  then have  $y^2 * (x^{(n+2)}) \leq x^2 * (y^{(n+2)})$  for  $n$ 
  by (metis (full-types) ab-semigroup-mult-class.mult-ac(1) mult.commute
power-add)
  then have  $y^2 * (x^{(n+2)} / \text{fact } (n+2)) \leq x^2 * (y^{(n+2)} / \text{fact } (n+2))$ 
/ fact (n+2) for  $n$ 
  by (meson divide-right-mono fact-ge-zero)
  then have  $(\sum n. y^2 * (x^{(n+2)} / \text{fact } (n+2))) \leq (\sum n. x^2 * (y^{(n+2)} / \text{fact } (n+2)))$ 
apply (intro suminf-le[OF - ls rs])
by auto
  then have  $y^2 * (\exp x - x - 1) \leq x^2 * (\exp y - y - 1)$ 
using l r by presburger
}
moreover {
  assume ineq:  $|y| < |x|$ 

  from psi-mono[OF assms(1)]
  have  $(\exp x - x - 1) - x^2 / 2 \leq (\exp y - y - 1) - y^2 / 2$  .

  then have  $y^2 * ((\exp x - x - 1) - x^2 / 2) \leq x^2 * ((\exp y - y - 1) - y^2 / 2)$ 
  by (smt (verit, best) ineq diff-divide-distrib exp-lower-Taylor-quadratic
le(1) le(2) mult-nonneg-nonneg one-less-exp-iff power-zero-numeral prob-space.psi-mono
prob-space-completion right-diff-distrib zero-le-power2))

  then have  $y^2 * (\exp x - x - 1) \leq x^2 * (\exp y - y - 1)$ 
  by (simp add: mult.commute right-diff-distrib)
}
ultimately show ?thesis by auto
qed

```

```

lemma bennett-inequality-1:
  assumes I: finite I
  assumes ind: indep-vars ( $\lambda . \text{borel}$ ) X I
  assumes intsq:  $\bigwedge i. i \in I \implies \text{integrable } M (\lambda x. (X i x)^2)$ 

```

```

assumes bnd:  $\bigwedge i. i \in I \implies AE\ x\ in\ M. X\ i\ x \leq 1$ 
assumes t:  $t \geq 0$ 
defines V  $\equiv (\sum i \in I. expectation(\lambda x. X\ i\ x^2))$ 
shows prob { $x \in space\ M. (\sum i \in I. X\ i\ x - expectation\ (X\ i)) \geq$ 
t}  $\leq$ 
  exp (-V * bennett-h (t / V))
proof (cases V = 0)
  case True
  then show ?thesis
  by auto
next
  case f: False
  have V  $\geq 0$ 
  unfolding V-def
  apply (intro sum-nonneg integral-nonneg-AE)
  by auto
  then have Vpos: V > 0 using f by auto

define l :: real where l = ln(1 + t / V)
then have l: l  $\geq 0$ 
  using t Vpos by auto
have rv[measurable]:  $\bigwedge i. i \in I \implies random\ variable\ borel\ (X\ i)$ 
  using ind unfolding indep-vars-def by blast

define  $\psi$  where  $\psi = (\lambda x::real. exp(x) - x - 1)$ 

have rw: exp y = 1 + y +  $\psi$  y for y
  unfolding  $\psi$ -def by auto

have ebnd:  $\bigwedge i. i \in I \implies$ 
  AE x in M. exp (l * X i x)  $\leq exp\ l$ 
  apply (drule bnd)
  using l by (auto simp add: mult-left-le)

have int:  $\bigwedge i. i \in I \implies integrable\ M\ (\lambda x. (X\ i\ x))$ 
using rv intsq square-integrable-imp-integrable by blast

have intl:  $\bigwedge i. i \in I \implies integrable\ M\ (\lambda x. (l * X\ i\ x))$ 
using int by blast

have interpl:  $\bigwedge i. i \in I \implies integrable\ M\ (\lambda x. exp\ (l * X\ i\ x))$ 
apply (intro integrable-const-bound[where B = exp l])
using ebnd by auto

have intpsi:  $\bigwedge i. i \in I \implies integrable\ M\ (\lambda x. \psi\ (l * X\ i\ x))$ 
unfolding  $\psi$ -def
using intl interpl by auto

```

have **: $\bigwedge i. i \in I \implies$
 $expectation (\lambda x. \psi (l * X i x)) \leq \psi l * expectation (\lambda x. (X i x)^{\wedge 2})$
proof –
fix i **assume** $i: i \in I$
then have $AE\ x\ in\ M. l * X\ i\ x \leq l$
using $ebnd$ **by** $auto$
then have $AE\ x\ in\ M. l^{\wedge 2} * \psi (l * X i x) \leq (l * X i x)^{\wedge 2} * \psi l$
using $psi-inequality[OF\ -\ l]$ **unfolding** $\psi-def$
by $auto$
then have $AE\ x\ in\ M. l^{\wedge 2} * \psi (l * X i x) \leq l^{\wedge 2} * (\psi l * (X\ i\ x)^{\wedge 2})$
by $(auto\ simp\ add: field-simps)$
then have $AE\ x\ in\ M. \psi (l * X i x) \leq \psi l * (X i x)^{\wedge 2}$
by $(smt\ (verit,\ best)\ AE-cong\ \psi-def\ exp-eq-one-iff\ mult-cancel-left\ mult-eq-0-iff\ mult-left-mono\ zero-eq-power2\ zero-le-power2)$
then have $AE\ x\ in\ M. 0 \leq \psi l * (X i x)^{\wedge 2} - \psi (l * X i x)$
by $auto$
then have $expectation (\lambda x. \psi l * (X i x)^{\wedge 2} + (- \psi (l * X i x)))$
 ≥ 0
by $(simp\ add: integral-nonneg-AE)$
also have $expectation (\lambda x. \psi l * (X i x)^{\wedge 2} + (- \psi (l * X i x))) =$
 $\psi l * expectation (\lambda x. (X i x)^{\wedge 2}) - expectation (\lambda x. \psi (l * X i$
 $x))$
apply $(subst\ Bochner-Integration.integral-add)$
using $intpsi[OF\ i]$ $intsq[OF\ i]$ **by** $auto$
finally show $expectation (\lambda x. \psi (l * X i x)) \leq \psi l * expectation$
 $(\lambda x. (X i x)^{\wedge 2})$
by $auto$
qed

then have *: $\bigwedge i. i \in I \implies$
 $expectation (\lambda x. exp (l * X i x)) \leq$
 $exp (l * expectation (X i)) * exp (\psi l * expectation (\lambda x. X i x^{\wedge 2}))$
proof –
fix i
assume $iI: i \in I$
have $expectation (\lambda x. exp (l * X i x)) =$
 $1 + l * expectation (\lambda x. X i x) +$
 $expectation (\lambda x. \psi (l * X i x))$
unfolding rw
apply $(subst\ Bochner-Integration.integral-add)$
using $iI\ intl\ intpsi$ **apply** $auto[2]$
apply $(subst\ Bochner-Integration.integral-add)$
using $intl\ iI\ prob-space$ **by** $auto$
also have $\dots = l * expectation (X i) + 1 + expectation (\lambda x. \psi (l$
 $* X i x))$
by $auto$
also have $\dots \leq 1 + l * expectation (X i) + \psi l * expectation (\lambda x.$

$X i x \hat{2}$)
using $**[OF iI]$ **by** *auto*
also have $\dots \leq \exp (l * \text{expectation } (X i)) * \exp (\psi l * \text{expectation } (\lambda x. X i x \hat{2}))$
by (*simp add: is-num-normalize(1) mult-exp-exp*)
finally show $\text{expectation } (\lambda x. \exp (l * X i x)) \leq \exp (l * \text{expectation } (X i)) * \exp (\psi l * \text{expectation } (\lambda x. X i x \hat{2}))$

qed

have $(\prod_{i \in I. \text{expectation } (\lambda x. \exp (l * (X i x)))}) \leq (\prod_{i \in I. \exp (l * \text{expectation } (X i)) * \exp (\psi l * \text{expectation } (\lambda x. X i x \hat{2}))$)
by (*auto intro!: prod-mono simp add: **)
also have $\dots = (\prod_{i \in I. \exp (l * \text{expectation } (X i))}) * (\prod_{i \in I. \exp (\psi l * \text{expectation } (\lambda x. X i x \hat{2}))$)
by (*auto simp add: prod.distrib*)
finally have $**:$
 $(\prod_{i \in I. \text{expectation } (\lambda x. \exp (l * (X i x)))}) \leq (\prod_{i \in I. \exp (l * \text{expectation } (X i))}) * \exp (\psi l * V)$
by (*simp add: V-def I exp-sum sum-distrib-left*)

from *indep-vars-Chernoff-ineq-ge[OF I ind l interpl]*
have $\text{prob } \{x \in \text{space } M. (\sum i \in I. X i x - \text{expectation } (X i)) \geq t\} \leq$
 $\exp (-l * t) * (\prod_{i \in I. \text{expectation } (\lambda x. \exp (l * (X i x - \text{expectation } (X i))))}$
by *auto*
also have $(\prod_{i \in I. \text{expectation } (\lambda x. \exp (l * (X i x - \text{expectation } (X i))))}) = (\prod_{i \in I. \text{expectation } (\lambda x. \exp (l * (X i x)))} * \exp (-l * \text{expectation } (X i)))$
by (*auto intro!: prod.cong simp add: field-simps exp-diff exp-minus-inverse*)
also have $\dots = (\prod_{i \in I. \exp (-l * \text{expectation } (X i))}) * (\prod_{i \in I. \text{expectation } (\lambda x. \exp (l * (X i x)))}$)
by (*auto simp add: prod.distrib*)
also have $\dots \leq (\prod_{i \in I. \exp (-l * \text{expectation } (X i))}) * ((\prod_{i \in I. \exp (l * \text{expectation } (X i))}) * \exp (\psi l * V))$
apply (*intro mult-left-mono[OF **]*)
by (*meson exp-ge-zero prod-nonneg*)
also have $\dots = \exp (\psi l * V)$
apply (*simp add: prod.distrib [symmetric]*)
by (*smt (verit, ccfv-threshold) exp-minus-inverse prod.not-neutral-contains-not-neutral*)
finally have
 $\text{prob } \{x \in \text{space } M. (\sum i \in I. X i x - \text{expectation } (X i)) \geq t\} \leq \exp (\psi l * V - l * t)$

by (*simp add: mult-exp-exp*)
 also have $\psi \ l * V - l * t = -V * \text{bennett-h } (t / V)$
 unfolding $\psi\text{-def } l\text{-def } \text{bennett-h-def}$
 apply (*subst exp-ln*)
 subgoal by (*smt (verit) Vpos divide-nonneg-nonneg t*)
 by (*auto simp add: algebra-simps*)
 finally show *?thesis* .
 qed

lemma *real-AE-le-sum:*

assumes $\bigwedge i. i \in I \implies AE\ x\ in\ M. f\ i\ x \leq (g\ i\ x::real)$
 shows $AE\ x\ in\ M. (\sum i \in I. f\ i\ x) \leq (\sum i \in I. g\ i\ x)$
 proof (*cases*)
 assume *finite I*
 with *AE-finite-allI [OF this assms]* have $0: AE\ x\ in\ M. (\forall i \in I. f\ i\ x \leq g\ i\ x)$ by *auto*
 show *?thesis* by (*intro eventually-mono [OF 0] sum-mono*) *auto*
 qed *simp*

lemma *real-AE-eq-sum:*

assumes $\bigwedge i. i \in I \implies AE\ x\ in\ M. f\ i\ x = (g\ i\ x::real)$
 shows $AE\ x\ in\ M. (\sum i \in I. f\ i\ x) = (\sum i \in I. g\ i\ x)$
 proof –
 have $1: AE\ x\ in\ M. (\sum i \in I. f\ i\ x) \leq (\sum i \in I. g\ i\ x)$
 apply (*intro real-AE-le-sum*)
 apply (*drule assms*)
 by *auto*
 have $2: AE\ x\ in\ M. (\sum i \in I. g\ i\ x) \leq (\sum i \in I. f\ i\ x)$
 apply (*intro real-AE-le-sum*)
 apply (*drule assms*)
 by *auto*
 show *?thesis*
 using $1\ 2$
 by *auto*
 qed

theorem *bennett-inequality:*

assumes *I: finite I*
 assumes *ind: indep-vars* ($\lambda \cdot. \text{borel}$) $X\ I$
 assumes *intsq:* $\bigwedge i. i \in I \implies \text{integrable } M\ (\lambda x. (X\ i\ x)^{\wedge 2})$
 assumes *bnd:* $\bigwedge i. i \in I \implies AE\ x\ in\ M. X\ i\ x \leq B$
 assumes *t:* $t \geq 0$
 assumes *B:* $B > 0$
 defines $V \equiv (\sum i \in I. \text{expectation } (\lambda x. X\ i\ x^{\wedge 2}))$
 shows $\text{prob } \{x \in \text{space } M. (\sum i \in I. X\ i\ x - \text{expectation } (X\ i)) \geq t\} \leq$
 $\text{exp } (-V / B^{\wedge 2} * \text{bennett-h } (t * B / V))$
 proof –

```

define  $Y$  where  $Y = (\lambda i x. X i x / B)$ 

from indep-vars-compose[OF ind, where  $Y = \lambda i x. x / B$ ]
have  $1$ : indep-vars ( $\lambda \cdot$ . borel)  $Y I$ 
  unfolding  $Y$ -def by (auto simp add: o-def)
have  $2$ :  $\bigwedge i. i \in I \implies \text{integrable } M (\lambda x. (Y i x)^2)$ 
  unfolding  $Y$ -def apply (drule intsq)
  by (auto simp add: field-simps)
have  $3$ :  $\bigwedge i. i \in I \implies AE x \text{ in } M. Y i x \leq 1$ 
  unfolding  $Y$ -def apply (drule bnd)
  using  $B$  by auto
have  $4$ :  $0 \leq t / B$  using  $t B$  by auto

have  $rw1$ :  $(\sum i \in I. Y i x - \text{expectation } (Y i)) =$ 
 $(\sum i \in I. X i x - \text{expectation } (X i)) / B$  for  $x$ 
  unfolding  $Y$ -def
  by (auto simp: diff-divide-distrib sum-divide-distrib)

have  $rw2$ :  $\text{expectation } (\lambda x. (Y i x)^2) =$ 
 $\text{expectation } (\lambda x. (X i x)^2) / B^2$  for  $i$ 
  unfolding  $Y$ -def
  by (simp add: power-divide)

have  $rw3$ :  $-(\sum i \in I. \text{expectation } (\lambda x. (X i x)^2) / B^2) = - V /$ 
 $B^2$ 
  unfolding  $V$ -def
  by (auto simp add: sum-divide-distrib)

have  $t / B / (\sum i \in I. \text{expectation } (\lambda x. (X i x)^2) / B^2) =$ 
 $t / B / (V / B^2)$ 
  unfolding  $V$ -def
  by (auto simp add: sum-divide-distrib)
then have  $rw4$ :  $t / B / (\sum i \in I. \text{expectation } (\lambda x. (X i x)^2) / B^2)$ 
 $=$ 
 $t * B / V$ 
  by (simp add: power2-eq-square)
have  $\text{prob } \{x \in \text{space } M. t \leq (\sum i \in I. X i x - \text{expectation } (X i))\}$ 
 $=$ 
 $\text{prob}\{x \in \text{space } M. t / B \leq (\sum i \in I. X i x - \text{expectation } (X i)) /$ 
 $B\}$ 
  by (smt (verit, best) B Collect-cong divide-cancel-right divide-right-mono)
also have  $\dots \leq$ 
 $\text{exp } (- V / B^2 * \text{bennett-h } (t * B / V))$ 
  using bennett-inequality-1[OF I 1 2 3 4]
  unfolding  $rw1$   $rw2$   $rw3$   $rw4$  .
finally show ?thesis .
qed

```



```

lemma bennett-h-bernstein-bound:
  assumes  $x \geq 0$ 
  shows  $\text{bennett-h } x \geq x^2 / (2 * (1 + x / 3))$ 
proof -
  have  $\text{eq}: x^2 / (2 * (1 + x / 3)) = 3/2 * x - 9/2 * (x / (x+3))$ 
    using assms
    by (sos (()) & ())

  define g where  $g = (\lambda x. \text{bennett-h } x - (3/2 * x - 9/2 * (x / (x+3))))$ 

  define g' where  $g' = (\lambda x::\text{real}. \ln(1 + x) + 27 / (2 * (x+3)^2) - 3 / 2)$ 
  define g'' where  $g'' = (\lambda x::\text{real}. 1 / (1 + x) - 27 / (x+3)^3)$ 

  have  $54 / ((2 * x + 6)^2) = 27 / (2 * (x + 3)^2)$  (is ?L = ?R)
for  $x :: \text{real}$ 
proof -
  have  $?L = 54 / (2^2 * (x + 3)^2)$ 
  unfolding power-mult-distrib[symmetric] by (simp add: algebra-simps)
  also have  $\dots = ?R$  by simp
  finally show thesis by simp
qed

hence  $1: x \geq 0 \implies (g \text{ has-real-derivative } (g' x)) \text{ (at } x \text{) for } x$ 
  unfolding g-def g'-def bennett-h-def by (auto intro!: derivative-eq-intros simp: power2-eq-square)
have  $2: x \geq 0 \implies (g' \text{ has-real-derivative } (g'' x)) \text{ (at } x \text{) for } x$ 
  unfolding g'-def g''-def
  apply (auto intro!: derivative-eq-intros)[1]
  by (sos (()) & ())

have gz:  $g 0 = 0$ 
  unfolding g-def bennett-h-def by auto
have g1z:  $g' 0 = 0$ 
  unfolding g'-def by auto

have p2:  $g'' x \geq 0$  if  $x \geq 0$  for  $x$ 
proof -
  have  $27 * (1+x) \leq (x+3)^3$ 
  using that unfolding power3-eq-cube by (auto simp: algebra-simps)
  hence  $27 / (x + 3)^3 \leq 1 / (1+x)$ 
  using that by (subst frac-le-eq) (auto intro!: divide-nonpos-pos)
  thus thesis unfolding g''-def by simp
qed

from deriv-nonneg-imp-mono[OF 2 p2 -]

```

have $x \geq 0 \implies g' x \geq 0$ **for** x **using** $g1z$
by (*metis atLeastAtMost-iff*)

from *deriv-nonneg-imp-mono*[*OF 1 this -*]
have $x \geq 0 \implies g x \geq 0$ **for** x **using** gz
by (*metis atLeastAtMost-iff*)

thus *?thesis*
using *assms eq g-def* **by** *force*
qed

lemma *sum-sq-exp-eq-zero-imp-zero*:
assumes *finite I i ∈ I*
assumes *intsq: integrable M (λx. (X i x)^2)*
assumes $(\sum i \in I. \text{expectation } (\lambda x. X i x^2)) = 0$
shows *AE x in M. X i x = (0::real)*
proof –
have $(\forall i \in I. \text{expectation } (\lambda x. X i x^2) = 0)$
using *assms*
apply (*subst sum-nonneg-eq-0-iff[symmetric]*)
by *auto*
then have $\text{expectation } (\lambda x. X i x^2) = 0$
using *assms(2)* **by** *blast*
thus *?thesis*
using *integral-nonneg-eq-0-iff-AE*[*OF intsq*]
by *auto*
qed

corollary *bernstein-inequality*:
assumes *I: finite I*
assumes *ind: indep-vars (λ -, borel) X I*
assumes *intsq: $\bigwedge i. i \in I \implies \text{integrable } M (\lambda x. (X i x)^2)$*
assumes *bnf: $\bigwedge i. i \in I \implies \text{AE } x \text{ in } M. X i x \leq B$*
assumes *t: $t \geq 0$*
assumes *B: $B > 0$*
defines $V \equiv (\sum i \in I. \text{expectation } (\lambda x. X i x^2))$
shows $\text{prob } \{x \in \text{space } M. (\sum i \in I. X i x - \text{expectation } (X i)) \geq t\} \leq$
 $\text{exp } (- (t^2 / (2 * (V + t * B / 3))))$
proof (*cases V = 0*)
case *True*
then have $1: \bigwedge i. i \in I \implies \text{AE } x \text{ in } M. X i x = 0$
unfolding *V-def*
using *sum-sq-exp-eq-zero-imp-zero*
by (*metis I intsq*)
then have $2: \bigwedge i. i \in I \implies \text{expectation } (X i) = 0$
using *integral-eq-zero-AE* **by** *blast*

have *AE x in M. $(\sum i \in I. X i x - \text{expectation } (X i)) = (\sum i \in I.$*

0)

```

apply (intro real-AE-eq-sum)
using 1 2
by auto
then have *: AE x in M. ( $\sum i \in I. X i x - \text{expectation } (X i) = 0$ )
by force

moreover {
  assume t > 0
  then have prob {x  $\in$  space M. ( $\sum i \in I. X i x - \text{expectation } (X$ 
i)  $\geq$  t} = 0
  apply (intro prob-eq-0-AE)
  using * by auto
  then have ?thesis by auto
}
ultimately show ?thesis
apply (cases t = 0) using t by auto
next
case f: False
have V  $\geq$  0
  unfolding V-def
  apply (intro sum-nonneg integral-nonneg-AE)
  by auto
then have V: V > 0 using f by auto

have t * B / V  $\geq$  0 using t B V by auto
from bennett-h-bernstein-bound[OF this]
have (t * B / V)2 / (2 * (1 + t * B / V / 3))
   $\leq$  bennett-h (t * B / V) .

then have (- V / B2) * bennett-h (t * B / V)  $\leq$ 
  (- V / B2) * ((t * B / V)2 / (2 * (1 + t * B / V / 3)))
  apply (subst mult-left-mono-neg)
  using B V by auto
also have ... =
  ((- V / B2) * (t * B / V)2) / (2 * (1 + t * B / V / 3))
  by auto
also have ((- V / B2) * (t * B / V)2) = -(t2) / V
  using V B by (auto simp add: field-simps power2-eq-square)
finally have *: (- V / B2) * bennett-h (t * B / V)  $\leq$ 
  -(t2) / (2 * (V + t * B / 3))
  using V by (auto simp add: field-simps)

from bennett-inequality[OF assms(1-6)]
have prob {x  $\in$  space M. ( $\sum i \in I. X i x - \text{expectation } (X i) \geq$  t}
 $\leq$ 
  exp (- V / B2 * bennett-h (t * B / V))
  using V-def by auto
also have ...  $\leq$  exp (- (t2 / (2 * (V + t * B / 3))))

```

```

    using *
    by auto
    finally show ?thesis .
qed

end

end

```

3 Bienaymé's identity

Bienaymé's identity [5, §17] can be used to deduce the variance of a sum of random variables, if their co-variance is known. A common use-case of the identity is the computation of the variance of the mean of pair-wise independent variables.

theory *Bienaymes-Identity*

```

imports Concentration-Inequalities-Preliminary
begin

```

```

context prob-space
begin

```

lemma *variance-divide:*

```

fixes  $f :: 'a \Rightarrow \text{real}$ 
assumes integrable M f
shows  $\text{variance } (\lambda\omega. f \ \omega / r) = \text{variance } f / r^2$ 
using assms
by (subst Bochner-Integration.integral-divide[OF assms(1)])
    (simp add:diff-divide-distrib[symmetric] power2-eq-square algebra-simps)

```

definition *covariance where*

```

covariance f g = expectation  $(\lambda\omega. (f \ \omega - \text{expectation } f) * (g \ \omega - \text{expectation } g))$ 

```

lemma *covariance-eq:*

```

fixes  $f :: 'a \Rightarrow \text{real}$ 
assumes  $f \in \text{borel-measurable } M$   $g \in \text{borel-measurable } M$ 
assumes integrable M ( $\lambda\omega. f \ \omega^2$ ) integrable M ( $\lambda\omega. g \ \omega^2$ )
shows  $\text{covariance } f g = \text{expectation } (\lambda\omega. f \ \omega * g \ \omega) - \text{expectation } f$ 
     $* \text{expectation } g$ 
proof –
have integrable M f using square-integrable-imp-integrable assms by
auto
moreover have integrable M g using square-integrable-imp-integrable
assms by auto
ultimately show ?thesis
    using assms cauchy-schwartz(1)[where M=M]
    by (simp add:covariance-def algebra-simps prob-space)

```

qed

lemma *covar-integrable*:

fixes $f\ g :: 'a \Rightarrow \text{real}$

assumes $f \in \text{borel-measurable } M\ g \in \text{borel-measurable } M$

assumes $\text{integrable } M (\lambda\omega. f\ \omega^{\wedge}2)\ \text{integrable } M (\lambda\omega. g\ \omega^{\wedge}2)$

shows $\text{integrable } M (\lambda\omega. (f\ \omega - \text{expectation } f) * (g\ \omega - \text{expectation } g))$

proof –

have $\text{integrable } M\ f$ **using** *square-integrable-imp-integrable assms by auto*

moreover **have** $\text{integrable } M\ g$ **using** *square-integrable-imp-integrable assms by auto*

ultimately show *?thesis using assms cauchy-schwartz(1)[where $M=M$]* **by** (*simp add: algebra-simps*)

qed

lemma *sum-square-int*:

fixes $f :: 'b \Rightarrow 'a \Rightarrow \text{real}$

assumes *finite I*

assumes $\bigwedge i. i \in I \implies f\ i \in \text{borel-measurable } M$

assumes $\bigwedge i. i \in I \implies \text{integrable } M (\lambda\omega. f\ i\ \omega^{\wedge}2)$

shows $\text{integrable } M (\lambda\omega. (\sum i \in I. f\ i\ \omega)^2)$

proof –

have $\text{integrable } M (\lambda\omega. \sum i \in I. \sum j \in I. f\ j\ \omega * f\ i\ \omega)$

using *assms*

by (*intro Bochner-Integration.integrable-sum cauchy-schwartz(1)[where $M=M$], auto*)

thus *?thesis*

by (*simp add: power2-eq-square sum-distrib-left sum-distrib-right*)

qed

theorem *bienaymes-identity*:

fixes $f :: 'b \Rightarrow 'a \Rightarrow \text{real}$

assumes *finite I*

assumes $\bigwedge i. i \in I \implies f\ i \in \text{borel-measurable } M$

assumes $\bigwedge i. i \in I \implies \text{integrable } M (\lambda\omega. f\ i\ \omega^{\wedge}2)$

shows

$\text{variance } (\lambda\omega. (\sum i \in I. f\ i\ \omega)) = (\sum i \in I. (\sum j \in I. \text{covariance } (f\ i)\ (f\ j)))$

proof –

have $a: \bigwedge i\ j. i \in I \implies j \in I \implies$

$\text{integrable } M (\lambda\omega. (f\ i\ \omega - \text{expectation } (f\ i)) * (f\ j\ \omega - \text{expectation } (f\ j)))$

using *assms covar-integrable by simp*

have $\text{variance } (\lambda\omega. (\sum i \in I. f\ i\ \omega)) = \text{expectation } (\lambda\omega. (\sum i \in I. f\ i\ \omega - \text{expectation } (f\ i))^2)$

using *square-integrable-imp-integrable[OF assms(2,3)]*

by (*simp add: Bochner-Integration.integral-sum sum-subtractf*)

also have ... = expectation ($\lambda\omega. (\sum i \in I. (\sum j \in I. (f i \omega - \text{expectation } (f i)) * (f j \omega - \text{expectation } (f j))))$)
by (simp add: power2-eq-square sum-distrib-right sum-distrib-left mult.commute)
also have ... = ($\sum i \in I. (\sum j \in I. \text{covariance } (f i) (f j))$)
using a **by** (simp add: Bochner-Integration.integral-sum covariance-def)
finally show ?thesis **by** simp
qed

lemma covar-self-eq:
fixes f :: 'a \Rightarrow real
shows covariance f f = variance f
by (simp add: covariance-def power2-eq-square)

lemma covar-indep-eq-zero:
fixes f g :: 'a \Rightarrow real
assumes integrable M f
assumes integrable M g
assumes indep-var borel f borel g
shows covariance f g = 0

proof –
have a: indep-var borel (($\lambda t. t - \text{expectation } f$) \circ f) borel (($\lambda t. t - \text{expectation } g$) \circ g)
by (rule indep-var-compose[OF assms(3)], auto)

have b: expectation ($\lambda\omega. (f \omega - \text{expectation } f) * (g \omega - \text{expectation } g)$) = 0
using a **assms** **by** (subst indep-var-lebesgue-integral, auto simp add: comp-def prob-space)

thus ?thesis **by** (simp add: covariance-def)
qed

lemma bienaymes-identity-2:
fixes f :: 'b \Rightarrow 'a \Rightarrow real
assumes finite I
assumes $\bigwedge i. i \in I \implies f i \in \text{borel-measurable } M$
assumes $\bigwedge i. i \in I \implies \text{integrable } M (\lambda\omega. f i \omega^2)$
shows variance ($\lambda\omega. (\sum i \in I. f i \omega)$) =
 $(\sum i \in I. \text{variance } (f i)) + (\sum i \in I. \sum j \in I - \{i\}. \text{covariance } (f i) (f j))$

proof –
have variance ($\lambda\omega. (\sum i \in I. f i \omega)$) = ($\sum i \in I. \sum j \in I. \text{covariance } (f i) (f j)$)
by (simp add: bienaymes-identity[OF assms(1,2,3)])
also have ... = ($\sum i \in I. \text{covariance } (f i) (f i) + (\sum j \in I - \{i\}. \text{covariance } (f i) (f j))$)
using assms **by** (subst sum.insert[symmetric], auto simp add: insert-absorb)

also have ... = $(\sum_{i \in I}. \text{variance } (f \ i)) + (\sum_{i \in I}. (\sum_{j \in I - \{i\}}. \text{covariance } (f \ i) \ (f \ j)))$
by (*simp add: covar-self-eq[symmetric] sum.distrib*)
finally show ?thesis **by simp**
qed

theorem *bienaymes-identity-pairwise-indep:*

fixes $f :: 'b \Rightarrow 'a \Rightarrow \text{real}$
assumes *finite I*
assumes $\bigwedge i. i \in I \Rightarrow f \ i \in \text{borel-measurable } M$
assumes $\bigwedge i. i \in I \Rightarrow \text{integrable } M \ (\lambda \omega. f \ i \ \omega^{\wedge 2})$
assumes $\bigwedge i \ j. i \in I \Rightarrow j \in I \Rightarrow i \neq j \Rightarrow \text{indep-var borel } (f \ i)$
borel (f j)
shows $\text{variance } (\lambda \omega. (\sum_{i \in I}. f \ i \ \omega)) = (\sum_{i \in I}. \text{variance } (f \ i))$
proof –
have $\bigwedge i \ j. i \in I \Rightarrow j \in I - \{i\} \Rightarrow \text{covariance } (f \ i) \ (f \ j) = 0$
using *covar-indep-eq-zero assms(4) square-integrable-imp-integrable[OF assms(2,3)] by auto*
hence $a: (\sum_{i \in I}. \sum_{j \in I - \{i\}}. \text{covariance } (f \ i) \ (f \ j)) = 0$
by simp
thus ?thesis **by** (*simp add: bienaymes-identity-2[OF assms(1,2,3)]*)
qed

lemma *indep-var-from-indep-vars:*

assumes $i \neq j$
assumes *indep-vars* $(\lambda-. M') \ f \ \{i, j\}$
shows *indep-var* $M' \ (f \ i) \ M' \ (f \ j)$
proof –
have $a: \text{inj } (\text{case-bool } i \ j)$ **using** *assms(1)*
by (*simp add: bool.case-eq-if inj-def*)
have $b: \text{range } (\text{case-bool } i \ j) = \{i, j\}$
by (*simp add: UNIV-bool insert-commute*)
have $c: \text{indep-vars } (\lambda-. M') \ f \ (\text{range } (\text{case-bool } i \ j))$ **using** *assms(2)*
b by simp

have $\text{True} = \text{indep-vars } (\lambda x. M') \ (\lambda x. f \ (\text{case-bool } i \ j \ x)) \ \text{UNIV}$
using *indep-vars-reindex[OF a c]*
by (*simp add: comp-def*)
also have ... = *indep-vars* $(\lambda x. \text{case-bool } M' \ M' \ x) \ (\lambda x. \text{case-bool } (f \ i) \ (f \ j) \ x) \ \text{UNIV}$
by (*rule indep-vars-cong, auto simp: bool.case-distrib bool.case-eq-if*)
also have ... = ?thesis
by (*simp add: indep-var-def*)
finally show ?thesis **by simp**
qed

lemma *bienaymes-identity-pairwise-indep-2:*

fixes $f :: 'b \Rightarrow 'a \Rightarrow \text{real}$
assumes *finite I*

```

assumes  $\bigwedge i. i \in I \implies f i \in \text{borel-measurable } M$ 
assumes  $\bigwedge i. i \in I \implies \text{integrable } M (\lambda \omega. f i \omega^2)$ 
assumes  $\bigwedge J. J \subseteq I \implies \text{card } J = 2 \implies \text{indep-vars } (\lambda -. \text{borel}) f J$ 
shows  $\text{variance } (\lambda \omega. (\sum i \in I. f i \omega)) = (\sum i \in I. \text{variance } (f i))$ 
using assms(4)
by (intro bienaymes-identity-pairwise-indep[OF assms(1,2,3)] indep-var-from-indep-vars, auto)

```

lemma *bienaymes-identity-full-indep:*

```

fixes  $f :: 'b \Rightarrow 'a \Rightarrow \text{real}$ 
assumes finite I
assumes  $\bigwedge i. i \in I \implies f i \in \text{borel-measurable } M$ 
assumes  $\bigwedge i. i \in I \implies \text{integrable } M (\lambda \omega. f i \omega^2)$ 
assumes indep-vars  $(\lambda -. \text{borel}) f I$ 
shows  $\text{variance } (\lambda \omega. (\sum i \in I. f i \omega)) = (\sum i \in I. \text{variance } (f i))$ 
by (intro bienaymes-identity-pairwise-indep-2[OF assms(1,2,3)] indep-vars-subset[OF assms(4)] auto)

```

end

end

4 Cantelli's Inequality

Cantelli's inequality [3] is an improvement of Chebyshev's inequality for one-sided tail bounds.

theory *Cantelli-Inequality*

imports *HOL-Probability.Probability*

begin

context *prob-space*

begin

lemma *cantelli-arith:*

```

assumes  $a > (0::\text{real})$ 
shows  $(V + (V / a)^2) / (a + (V / a)^2) = V / (a^2 + V)$  (is ?L = ?R)

```

proof –

```

have  $?L = ((V * a^2 + V^2) / a^2) / ((a^2 + V)^2 / a^2)$ 
using assms by (intro arg-cong2[where f=(/)] (simp-all add:field-simps power2-eq-square))

```

```

also have  $\dots = (V * a^2 + V^2) / (a^2 + V)^2$ 

```

```

using assms unfolding divide-divide-times-eq by simp

```

```

also have  $\dots = V * (a^2 + V) / (a^2 + V)^2$ 

```

```

by (intro arg-cong2[where f=(/)] (simp-all add: algebra-simps power2-eq-square))

```

```

also have  $\dots = ?R$  by (simp add:power2-eq-square)

```


finally show *?thesis* by simp
qed

theorem *cantelli-inequality*:

assumes *[measurable]*: random-variable borel Z

assumes *intZsq*: integrable M ($\lambda z. Z z^2$)

assumes $a: a > 0$

shows $\text{prob } \{z \in \text{space } M. Z z - \text{expectation } Z \geq a\} \leq$
 $\text{variance } Z / (a^2 + \text{variance } Z)$

proof –

define u where $u = \text{variance } Z / a$

have $u: u \geq 0$

unfolding *u-def*

by (*simp add: a divide-nonneg-pos*)

define Y where $Y = (\lambda z. Z z + (-\text{expectation } Z))$

have random-variable borel ($\lambda z. |Y z + u|$)

unfolding *Y-def*

by *auto*

then have $ev: \{z \in \text{space } M. a + u \leq |Y z + u|\} \in \text{events}$

by *auto*

have *intZ*: integrable M Z

apply (*subst square-integrable-imp-integrable[OF - intZsq]*)

by *auto*

then have $i1$: integrable M ($\lambda z. (Z z - \text{expectation } Z + u)^2$)

unfolding *power2-sum power2-diff* **using** *intZsq*

by *auto*

have *intY*: integrable M Y

unfolding *Y-def* **using** *intZ* **by** *auto*

have *intYsq*: integrable M ($\lambda z. Y z^2$)

unfolding *Y-def power2-sum* **using** *intZsq intZ* **by** *auto*

have *expectation* $Y = 0$

unfolding *Y-def*

apply (*subst Bochner-Integration.integral-add[OF intZ]*)

using *prob-space* **by** *auto*

then have *expectation* ($\lambda z. (Y z + u)^2$) =

expectation ($\lambda z. (Y z)^2$) + u^2

unfolding *power2-sum*

apply (*subst Bochner-Integration.integral-add[OF - -]*)

using *intY intYsq* **apply** *auto*[2]

apply (*subst Bochner-Integration.integral-add[OF - -]*)

using *intY intYsq* **apply** *auto*[2]

using *prob-space* **by** *auto*

then have $*$: *expectation* ($\lambda z. (Y z + u)^2$) = *variance* Z + u^2

unfolding *Y-def* **by** *auto*

have
 $\text{prob } \{z \in \text{space } M. Z z - \text{expectation } Z \geq a\} =$
 $\text{prob } \{z \in \text{space } M. Y z + u \geq a + u\}$
apply (*intro arg-cong*[**where** $f = \text{prob}$])
using *Y-def* **by** *auto*
also have $\dots \leq \text{prob } \{z \in \text{space } M. a + u \leq |Y z + u|\}$
apply (*intro finite-measure-mono*[*OF - ev*])
by *auto*

also have $\dots \leq \text{expectation } (\lambda z. (Y z + u)^2) / (a + u)^2$
apply (*intro second-moment-method*)
unfolding *Y-def* **using** *a u i1* **by** *auto*
also have $\dots = ((\text{variance } Z) + u^2) / (a + u)^2$
using *** **by** *auto*
also have $\dots = \text{variance } Z / (a^2 + \text{variance } Z)$
unfolding *u-def* **using** *a* **by** (*auto intro!*: *cantelli-arith*)
finally show *?thesis* .
qed

corollary *cantelli-inequality-neg*:
assumes [*measurable*]: *random-variable borel* Z
assumes *intZsq*: *integrable* $M (\lambda z. Z z^2)$
assumes $a: a > 0$
shows $\text{prob } \{z \in \text{space } M. Z z - \text{expectation } Z \leq -a\} \leq$
 $\text{variance } Z / (a^2 + \text{variance } Z)$
proof –
define nZ **where** [*simp*]: $nZ = (\lambda z. -Z z)$
have vnZ : *variance* $nZ = \text{variance } Z$
unfolding *nZ-def*
by (*auto simp add*: *power2-commute*)

have *1*: *random-variable borel* nZ **by** *auto*
have *2*: *integrable* $M (\lambda z. (nZ z)^2)$
using *intZsq* **by** *auto*
from *cantelli-inequality*[*OF 1 2 a*]
have $\text{prob } \{z \in \text{space } M. a \leq nZ z - \text{expectation } nZ\} \leq$
 $\text{variance } nZ / (a^2 + \text{variance } nZ)$
by *auto*
thus *?thesis* **unfolding** vnZ **apply** *auto*[*1*]
by (*smt* (*verit*, *del-insts*) *Collect-cong*)
qed

end

end

5 Efron-Stein Inequality

In this section we verify the Efron-Stein inequality. The verified theorem is stated as Efron-Stein inequality for non-symmetric functions by Steele [8]. However most textbook refer to this version as “the Efron-Stein inequality”. The original result that was shown by Efron and Stein is a tail bound for the variance of a symmetric functions of i.i.d. random variables [4].

theory *Efron-Stein-Inequality*

imports *Concentration-Inequalities-Preliminary*

begin

theorem *efron-stein-inequality-distr*:

fixes $f :: - \Rightarrow \text{real}$

assumes *finite I*

assumes $\bigwedge i. i \in I \implies \text{prob-space } (M\ i)$

assumes *integrable (PiM I M) (λx. f x^2) and f-meas: f ∈ borel-measurable (PiM I M)*

shows *prob-space.variance (PiM I M) f ≤*

$(\sum i \in I. (\int x. (f (\lambda j. x (j, \text{False})) - f (\lambda j. x (j, j=i)))^2 \partial \text{PiM } (I \times \text{UNIV}) (M \circ \text{fst}))) / 2$

(is ?L ≤ ?R)

proof –

let $?M = \text{PiM } (I \times (\text{UNIV} :: \text{bool set})) (M \circ \text{fst})$

have *prob: prob-space (PiM I M)*

using *assms(2) by (intro prob-space-PiM) auto*

interpret *prob-space ?M*

using *assms(2) by (intro prob-space-PiM) auto*

define n **where** $n = \text{card } I$

obtain $q :: - \Rightarrow \text{nat}$ **where** $q: \text{bij-betw } q\ I\ \{..<n\}$

unfolding *n-def using ex-bij-betw-finite-nat[OF assms(1)] atLeast0LessThan*
by *auto*

let $?φ = (\lambda n\ x. f (\lambda j. x (j, q\ j < n)))$

let $?τ = (\lambda n\ x. f (\lambda j. x (j, q\ j = n)))$

let $?σ = (\lambda x. f (\lambda j. x (j, \text{False})))$

let $?χ = (\lambda x. f (\lambda j. x (j, \text{True})))$

have *meas-1: (λω. f (g ω)) ∈ borel-measurable ?M*

if $g \in \text{PiM } (I \times \text{UNIV}) (M \circ \text{fst}) \rightarrow_M \text{PiM } I\ M$ **for** g

using *that by (intro measurable-compose[OF f-meas])*

have *meas-2: (λx j. x (j, h j)) ∈ ?M →M PiM I M for h*

proof –

have $?thesis \longleftrightarrow (\lambda x. (\lambda j \in I. x (j, h j))) \in ?M \rightarrow_M Pi_M I M$
by (intro measurable-cong) (auto simp:space-PiM PiE-def extensional-def)
also have ... $\longleftrightarrow True$
unfolding eq-True
by (intro measurable-restrict measurable-PiM-component-rev) auto
finally show $?thesis$ **by** simp
qed

have $int-1$: integrable $?M (\lambda x. (g x - h x)^2)$
if integrable $?M (\lambda x. (g x)^2)$ integrable $?M (\lambda x. (h x)^2)$
and $g \in \text{borel-measurable } ?M$ $h \in \text{borel-measurable } ?M$
for $g h :: - \Rightarrow real$
proof –
have integrable $?M (\lambda x. (g x)^2 + (h x)^2 - 2 * (g x * h x))$
using that **by** (intro Bochner-Integration.integrable-add Bochner-Integration.integrable-diff
integrable-mult-right cauchy-schwartz(1))
thus $?thesis$ **by** (simp add:algebra-simps power2-eq-square)
qed

note $meas-rules = \text{borel-measurable-add borel-measurable-times borel-measurable-diff}$
 $\text{borel-measurable-power meas-1 meas-2}$

have $f-int$: integrable $(Pi_M I M) f$
by (intro finite-measure.square-integrable-imp-integrable[OF - f-meas
assms(3)]
prob-space.finite-measure prob)
moreover have integrable $(Pi_M I M) (\lambda x. f (restrict x I)) = \text{integrable } (Pi_M I M) f$
by (intro Bochner-Integration.integrable-cong) (auto simp:space-PiM)
ultimately have $f-int-2$: integrable $(Pi_M I M) (\lambda x. f (restrict x I))$
by simp

have $cong$: $(\int x. g (\lambda j \in I. x (j, h j)) \partial ?M) = (\int x. g (\lambda j. x (j, h j)) \partial ?M)$ (is $?L1 = ?R1$)
for $g :: - \Rightarrow real$ **and** h
by (intro Bochner-Integration.integral-cong arg-cong[where $f=g$]
refl)
(auto simp add:space-PiM PiE-def extensional-def restrict-def)

have $lift$: $(\int x. g x \partial Pi_M I M) = (\int x. g (\lambda j. x (j, h j)) \partial ?M)$ (is $?L1 = ?R1$)
if $g \in \text{borel-measurable } (Pi_M I M)$
for $g :: - \Rightarrow real$ **and** h
proof –
let $?J = (\lambda i. (i, h i)) ' I$
have $?R1 = (\int x. g (\lambda j \in I. x (j, h j)) \partial ?M)$
by (intro cong[symmetric])
also have ... $= (\int x. g x \partial distr ?M (Pi_M I (\lambda i. (M \circ fst) (i, h i)))$

$(\lambda x. (\lambda j \in I. x (j, h j)))$
using that
by (*intro integral-distr[symmetric] measurable-restrict measurable-component-singleton*) *auto*
also have ... = $(\int x. g x \partial PiM I (\lambda i. (M \circ fst) (i, h i)))$
using *assms(2)*
by (*intro arg-cong2[where f=integral^L] refl distr-PiM-reindex inj-onI*) *auto*
also have ... = ?L1
by *auto*
finally show ?thesis
by *simp*
qed

have *lift-int: integrable ?M* $(\lambda x. g (\lambda j. x (j, h j)))$ **if** *integrable (PiM I M) g*
for $g :: - \Rightarrow \text{real}$ **and** h
proof –
have *0:integrable (distr ?M (PiM I (\lambda i. (M \circ fst) (i, h i)))* $(\lambda x. (\lambda j \in I. x (j, h j)))$ g
using that *assms(2)* **by** (*subst distr-PiM-reindex*) (*auto intro:inj-onI*)
have *integrable ?M* $(\lambda x. g (\lambda j \in I. x (j, h j)))$
by (*intro integrable-distr[OF - 0] measurable-restrict measurable-component-singleton*) *auto*
moreover have *integrable ?M* $(\lambda x. g (\lambda j \in I. x (j, h j))) \longleftrightarrow ?thesis$
by (*intro Bochner-Integration.integrable-cong refl arg-cong[where f=g] ext*)
(auto simp:PiE-def space-PiM extensional-def)
ultimately show ?thesis
by *simp*
qed

note *int-rules = cauchy-schwartz(1) int-1 lift-int assms(3) f-int f-int-2*

have $(\int x. g x \partial ?M) = (\int x. g (\lambda(j,v). x (j, v \neq h j)) \partial ?M)$ (**is** ?L1 = ?R1)
if $g :: - \Rightarrow \text{real}$ **and** h
proof –
have ?L1 = $(\int x. g x \partial \text{distr } ?M (PiM (I \times UNIV) (\lambda i. (M \circ fst) (fst i, snd i \neq h (fst i))))$
 $(\lambda x. (\lambda i \in I \times UNIV. x (fst i, snd i \neq h (fst i))))$
by (*subst distr-PiM-reindex*) (*auto intro:inj-onI assms(2) simp:comp-def*)
also have ... = $(\int x. g (\lambda i \in I \times UNIV. x (fst i, snd i \neq h (fst i))) \partial ?M)$
using that by (*intro integral-distr measurable-restrict measurable-component-singleton*)
(auto simp:comp-def)

also have ... = ?R1
by (intro Bochner-Integration.integral-cong refl arg-cong[where
f=g] ext)
(auto simp add:space-PiM PiE-def extensional-def restrict-def)
finally show ?thesis
by simp
qed

hence switch: $(\int x. g x \partial^?M) = (\int x. h x \partial^?M)$
if $\bigwedge x. h x = g (\lambda(j,v). x (j, v \neq u j))$ $g \in \text{borel-measurable } ?M$
for $g h :: - \Rightarrow \text{real}$ **and** u
using that **by** simp

have 1: $(\int x. (?\sigma x) * (?\varphi i x - ?\varphi (i+1) x) \partial^?M) \leq (\int x. (?\sigma x - ?\tau i x)^2 \partial^?M) / 2$
(is ?L1 \leq ?R1)
if $i < n$ **for** i
proof -
have ?L1 = $(\int x. (?\tau i x) * (?\varphi (i+1) x - ?\varphi i x) \partial^?M)$
by (intro switch[of - - ($\lambda j. q j = i$)] arg-cong2[where f=(*)]
arg-cong2[where f=(-)] arg-cong[where f=f] ext meas-rules)
(auto intro:arg-cong)
hence ?L1 = $(?L1 + (\int x. (?\tau i x) * (?\varphi (i+1) x - ?\varphi i x) \partial^?M)) / 2$
by simp
also have ... = $(\int x. (?\sigma x) * (?\varphi i x - ?\varphi(i+1) x) + (?\tau i x) * (?\varphi(i+1) x - ?\varphi i x) \partial^?M) / 2$
by (intro Bochner-Integration.integral-add[symmetric] arg-cong2[where
f=(/)] refl
int-rules meas-rules)
also have ... = $(\int x. (?\sigma x - ?\tau i x) * (?\varphi i x - ?\varphi(i+1) x) \partial^?M) / 2$
by (intro arg-cong2[where f=(/)] Bochner-Integration.integral-cong)
(auto simp:algebra-simps)
also have ... $\leq ((\int x. (?\sigma x - ?\tau i x)^2 \partial^?M)^{\text{powr}(1/2)} * (\int x. (?\varphi i x - ?\varphi(i+1) x)^2 \partial^?M)^{\text{powr}(1/2)}) / 2$
by (intro divide-right-mono cauchy-schwartz meas-rules int-rules)
auto
also have ... = $((\int x. (?\sigma x - ?\tau i x)^2 \partial^?M)^{\text{powr}(1/2)} * (\int x. (?\sigma x - ?\tau i x)^2 \partial^?M)^{\text{powr}(1/2)}) / 2$
by (intro arg-cong2[where f=(/)] arg-cong2[where f=(*)] arg-cong2[where
f=(powr)] refl
switch[of - - ($\lambda j. q j < i$)] arg-cong2[where f=powr] arg-cong2[where
f=(-)]
arg-cong[where f=f] ext meas-rules) (auto intro:arg-cong)
also have ... = $(\int x. (?\sigma x - ?\tau i x)^2 \partial^?M) / 2$
by (simp add:powr-add[symmetric])
finally show ?thesis **by** simp
qed

have *indep-vars* ($M \circ \text{fst}$) ($\lambda i \omega. \omega i$) ($I \times \text{UNIV}$)
using *assms(2)* **by** (*intro proj-indep*) *auto*
hence $2:\text{indep-var}$ ($Pi_M (I \times \{False\}) (M \circ \text{fst})$) ($\lambda x. \lambda j \in I \times \{False\}. x j$)
 $(Pi_M (I \times \{True\}) (M \circ \text{fst})) (\lambda x. \lambda j \in I \times \{True\}. x j)$
by (*intro indep-var-restrict*[**where** $I = I \times \text{UNIV}$]) *auto*
have *indep-var*
 $(Pi_M I M) ((\lambda x. (\lambda i \in I. x (i, False)))) \circ (\lambda x. (\lambda j \in I \times \{False\}. x j))$
 $(Pi_M I M) ((\lambda x. (\lambda i \in I. x (i, True)))) \circ (\lambda x. (\lambda j \in I \times \{True\}. x j))$
by (*intro indep-var-compose*[*OF 2*] *measurable-restrict measurable-PiM-component-rev*) *auto*
hence *indep-var* ($Pi_M I M$) ($\lambda x. (\lambda j \in I. x (j, False))$) ($Pi_M I M$)
 $(\lambda x. (\lambda j \in I. x (j, True)))$
unfolding *comp-def* **by** (*simp add:restrict-def cong:if-cong*)

hence *indep-var borel* ($f \circ (\lambda x. (\lambda j \in I. x (j, False)))$) *borel* ($f \circ (\lambda x. (\lambda j \in I. x (j, True)))$)
by (*intro indep-var-compose*[*OF - assms(4,4)*]) *auto*
hence *indep:indep-var borel* ($\lambda x. f (\lambda j \in I. x (j, False))$) *borel* ($\lambda x. f (\lambda j \in I. x (j, True))$)
by (*simp add:comp-def*)

have $3: \omega (j, q j = q i) = \omega (j, j = i)$ **if**
 $\omega \in PiE (I \times \text{UNIV}) (\lambda i. \text{space } (M (\text{fst } i))) i \in I$ **for** $i j \omega$
proof (*cases j \in I*)
case *True*
hence ($q j = q i$) = ($j = i$)
using *that inj-onD*[*OF bij-betw-imp-inj-on*[*OF q*]] **by** *blast*
thus *?thesis* **by** *simp*
next
case *False*
hence $\omega (j, a) = \text{undefined}$ **for** a
using *that unfolding PiE-def extensional-def* **by** *simp*
thus *?thesis* **by** *simp*
qed

have $?L = (\int x. (f x)^2 \partial PiM I M) - (\int x. (f x) \partial PiM I M)^2$
by (*intro prob-space.variance-eq f-int assms(3) prob*)
also have $\dots = (\int x. (f x)^2 \partial PiM I M) - (\int x. f x \partial PiM I M) * (\int x. f x \partial PiM I M)$
by (*simp add:power2-eq-square*)
also have $\dots = (\int x. (?\sigma x)^2 \partial ?M) - (\int x. ?\sigma x \partial ?M) * (\int x. ?\chi x \partial ?M)$
by (*intro arg-cong2*[**where** $f = (-)$] *lift arg-cong2*[**where** $f = (*)$])
meas-rules f-meas)
also have $\dots = (\int x. (?\sigma x)^2 \partial ?M) - (\int x. f (\lambda j \in I. x (j, False))$

$\partial^?M) * (\int x. f(\lambda j \in I. x(j, True)) \partial^?M)$
by (intro arg-cong2[**where** $f=(-)$] arg-cong2[**where** $f=(*)$] cong[symmetric] refl)
also have ... = $(\int x. (?\sigma x) \wedge^2 \partial^?M) - (\int x. f(\lambda j \in I. x(j, False)) * f(\lambda j \in I. x(j, True)) \partial^?M)$
by (intro arg-cong2[**where** $f=(-)$] indep-var-lebesgue-integral[symmetric] refl int-rules indep)
also have ... = $(\int x. (?\sigma x) * (?\varphi 0 x) \partial^?M) - (\int x. (?\sigma x) * (?\varphi n x) \partial^?M)$
using bij-betw-apply[OF q] **by** (intro arg-cong2[**where** $f=(-)$] arg-cong2[**where** $f=(*)$] ext
arg-cong[**where** $f=f$] Bochner-Integration.integral-cong)
(auto simp:space-PiM power2-eq-square PiE-def extensional-def)
also have ... = $(\sum i < n. (\int x. (?\sigma x) * (?\varphi i x) \partial^?M) - (\int x. (?\sigma x) * (?\varphi (Suc i) x) \partial^?M))$
unfolding power2-eq-square **by** (intro sum-lessThan-telescope'[symmetric])
also have ... = $(\sum i < n. (\int x. (?\sigma x) * (?\varphi i x) - (?\sigma x) * (?\varphi (Suc i) x) \partial^?M))$
by (intro sum.cong Bochner-Integration.integral-diff[symmetric] int-rules meas-rules) auto
also have ... = $(\sum i < n. (\int x. (?\sigma x) * (?\varphi i x - ?\varphi (i+1) x) \partial^?M))$
by (simp-all add:power2-eq-square algebra-simps)
also have ... $\leq (\sum i < n. ((\int x. (?\sigma x - ?\tau i x) \wedge^2 \partial^?M)) / 2)$
by (intro sum-mono 1) auto
also have ... = $(\sum i \in I. ((\int x. (f(\lambda j. x(j, False)) - f(\lambda j. x(j, q j=q i))) \wedge^2 \partial^?M)) / 2)$
by (intro sum.reindex-bij-betw[OF q, symmetric])
also have ... = $(\sum i \in I. ((\int x. (f(\lambda j. x(j, False)) - f(\lambda j. x(j, q j=q i))) \wedge^2 \partial^?M))) / 2$
unfolding sum-divide-distrib[symmetric] **by** simp
also have ... = ?R
using inj-onD[OF bij-betw-imp-inj-on[OF q]]
by (intro arg-cong2[**where** $f=(/)$] arg-cong2[**where** $f=(-)$] arg-cong2[**where** $f=power$] arg-cong[**where** $f=f$] Bochner-Integration.integral-cong sum.cong refl ext 3)
(auto simp add:space-PiM)
finally show ?thesis
by simp
qed

theorem (in prob-space) efron-stein-inequality-classic:

fixes $f :: - \Rightarrow real$

assumes finite I

assumes indep-vars ($M' \circ fst$) X (I \times (UNIV :: bool set))

assumes $f \in$ borel-measurable (PiM I M')

assumes integrable M ($\lambda \omega. f(\lambda i \in I. X(i, False) \omega) \wedge^2$)

assumes $\bigwedge i. i \in I \implies distr M (M' i) (X(i, True)) = distr M (M'$

i) ($X (i, False)$)
shows *variance* ($\lambda\omega. f (\lambda i \in I. X (i, False) \omega) \leq$
 $(\sum j \in I. expectation (\lambda\omega. (f (\lambda i \in I. X (i, False) \omega) - f (\lambda i \in I. X$
 $(i, i=j) \omega))^2) / 2$
 $(is ?L \leq ?R)$)

proof –
let $?D = distr M (PiM I M') (\lambda\omega. \lambda i \in I. X (i, False) \omega)$

let $?M = PiM I (\lambda i. distr M (M' i) (X (i, False)))$
let $?N = PiM (I \times (UNIV::bool set)) ((\lambda i. distr M (M' i) (X$
 $(i, False))) \circ fst)$

have *rv: random-variable* ($M' i$) ($X (i, j)$) **if** $i \in I$ **for** $i j$
using *assms(2)* that **unfolding** *indep-vars-def* **by** *auto*

have *proj-meas:* ($\lambda x j. x (j, h j) \in PiM (I \times UNIV) (M' \circ fst)$
 $\rightarrow_M PiM I M'$
for $h :: - \Rightarrow bool$

proof –
have $?thesis \longleftrightarrow (\lambda x. (\lambda j \in I. x (j, h j))) \in PiM (I \times UNIV)$
 $(M' \circ fst) \rightarrow_M PiM I M'$
by (*intro measurable-cong*) (*auto simp:space-PiM PiE-def exten-*
sional-def)
also have $\dots \longleftrightarrow True$
unfolding *eq-True*
by (*intro measurable-restrict measurable-PiM-component-rew*) *auto*
finally show $?thesis$ **by** *simp*

qed

note *meas-rules = borel-measurable-add borel-measurable-times borel-measurable-diff*
proj-meas
borel-measurable-power assms(3) measurable-restrict measurable-compose[OF
- assms(3)]

have *indep-vars* ($(M' \circ fst) \circ (\lambda i. (i, False))$) ($\lambda i. X (i, False)$) I
by (*intro indep-vars-reindex indep-vars-subset[OF assms(2)] inj-onI*)
auto
hence *indep-vars* $M' (\lambda i. X (i, False)) I$ **by** (*simp add: comp-def*)
hence $0: ?D = PiM I (\lambda i. distr M (M' i) (X (i, False)))$
by (*intro iffD1[OF indep-vars-iff-distr-eq-PiM'] rv*)

have $distr M (M' (fst x)) (X (fst x, False)) = distr M (M' (fst x))$
 $(X x)$
if $x \in I \times UNIV$ **for** x
using *that assms(5)* **by** (*cases x, cases snd x*) *auto*

hence $1: ?N = PiM (I \times UNIV) (\lambda i. distr M ((M' \circ fst) i) (X i))$
using *assms(3)* **by** (*intro PiM-cong refl*) (*simp add: comp-def*)
also have $\dots = distr M (PiM (I \times UNIV) (M' \circ fst)) (\lambda x. \lambda i \in I \times$

```

UNIV. X i x)
  using rv by (intro iffD1[OF indep-vars-iff-distr-eq-PiM'', symmetric]
  assms(2)) auto
  finally have 2: ?N = distr M (PiM (I × UNIV) (M' ∘ fst)) (λx.
  λi∈I × UNIV. X i x)
    by simp

  have 3: integrable (PiM I (λi. distr M (M' i) (X (i, False)))) (λx.
  (f x)2)
    unfolding 0[symmetric] by (intro iffD2[OF integrable-distr-eq]
  meas-rules assms rv)

  have ?L = (∫ x. (f x - expectation (λω. f (λi∈I. X (i, False) ω)))2
  ∂?D)
    using rv by (intro integral-distr[symmetric] meas-rules measur-
  able-restrict) auto
    also have ... = prob-space.variance ?D f
      by (intro arg-cong[where f=integralL ?D] arg-cong2[where f=(-)]
  arg-cong2[where f=power]
      refl ext integral-distr[symmetric] measurable-restrict rv assms(3))
    also have ... = prob-space.variance ?M f
      unfolding 0 by simp
    also have ... ≤ (∑ i∈I. (∫ x. (f (λj. x (j, False)) - f (λj. x (j, j =
  i)))2 ∂?N)) / 2
      using assms(3) by (intro efron-stein-inequality-distr prob-space-distr
  rv assms(1) 3) auto
    also have ... = (∑ i∈I. expectation (λω. (f (λj. (λi∈I × UNIV. X i
  ω) (j, False)) -
      f (λj. (λi∈I × UNIV. X i ω) (j, j=i))2))) / 2
      using rv unfolding 2
      by (intro sum.cong arg-cong2[where f=(/)] integral-distr refl
  meas-rules) auto
    also have ... = ?R
      by (simp add: restrict-def)
    finally show ?thesis
      by simp
qed

end

```

6 McDiarmid's inequality

In this section we verify McDiarmid's inequality [6, Lemma 1.2]. In the source and also further sources sometimes refer to the result as the “independent bounded differences” inequality.

```

theory McDiarmid-Inequality
  imports Concentration-Inequalities-Preliminary
begin

```

lemma *Collect-restr-cong*:
assumes $A = B$
assumes $\bigwedge x. x \in A \implies P x = Q x$
shows $\{x \in A. P x\} = \{x \in B. Q x\}$
using *assms* **by** *auto*

lemma *ineq-chain*:
fixes $h :: nat \Rightarrow real$
assumes $\bigwedge i. i < n \implies h (i+1) \leq h i$
shows $h n \leq h 0$
using *assms* **by** (*induction n*) *force+*

lemma *restrict-subset-eq*:
assumes $A \subseteq B$
assumes $restrict f B = restrict g B$
shows $restrict f A = restrict g A$
using *assms* **unfolding** *restrict-def* **by** (*meson subsetD*)

Bochner Integral version of Hoeffding's Lemma using *interval-bounded-random-variable.Hoeffding*

lemma (*in prob-space*) *Hoeffdings-lemma-bochner*:
assumes $l > 0$ **and** *E0*: *expectation f = 0*
assumes *random-variable borel f*
assumes *AE x in M. f x ∈ {a..b::real}*
shows $expectation (\lambda x. exp (l * f x)) \leq exp (l^2 * (b - a)^2 / 8)$ (*is ?L ≤ ?R*)

proof –

interpret *interval-bounded-random-variable M f a b*
using *assms* **by** (*unfold-locales*) *auto*

have *integrable M (λx. exp (l * f x))*
using *assms(1,3,4)* **by** (*intro integrable-const-bound[where B=exp (l * b)] simp-all*)

hence *ennreal (?L) = (∫⁺ x. exp (l * f x) ∂M)*
by (*intro nn-integral-eq-integral[symmetric]*) *auto*

also have $... \leq ennreal (?R)$

by (*intro Hoeffdings-lemma-nn-integral-0 assms*)

finally have $0:ennreal (?L) \leq ennreal ?R$

by *simp*

show *?thesis*

proof (*cases ?L ≥ 0*)

case *True*

thus *?thesis* **using** *0* **by** *simp*

next

case *False*

hence $?L \leq 0$ **by** *simp*

also have $... \leq ?R$ **by** *simp*

finally show *?thesis* **by** *simp*

qed
qed

lemma (in prob-space) Hoeffdings-lemma-bochner-2:

assumes $l > 0$ and $E0$: expectation $f = 0$

assumes random-variable borel f

assumes $\bigwedge x y. \{x,y\} \subseteq \text{space } M \implies |f x - f y| \leq (c::\text{real})$

shows expectation $(\lambda x. \exp (l * f x)) \leq \exp (l^2 * c^2 / 8)$ (is ?L
 $\leq ?R$)

proof -

define $a :: \text{real}$ where $a = (\text{INF } x \in \text{space } M. f x)$

define $b :: \text{real}$ where $b = a + c$

obtain ω where $\omega : \omega \in \text{space } M$ using not-empty by auto

hence $0 : f ' \text{space } M \neq \{\}$ by auto

have $1 : c = b - a$ unfolding b-def by simp

have bdd-below $(f ' \text{space } M)$

using ω assms(4) unfolding abs-le-iff

by (intro bdd-belowI[where $m=f \omega - c$]) (auto simp add:algebra-simps)

hence $f x \geq a$ if $x \in \text{space } M$ for x unfolding a-def by (intro
 $c\text{INF-lower that}$)

moreover have $f x \leq b$ if $x\text{-space}: x \in \text{space } M$ for x

proof (rule ccontr)

assume $\neg(f x \leq b)$

hence $a : f x > a + c$ unfolding b-def by simp

have $f y \geq f x - c$ if $y \in \text{space } M$ for y

using that $x\text{-space assms}(4)$ unfolding abs-le-iff by (simp
 add:algebra-simps)

hence $f x - c \leq a$ unfolding a-def using $c\text{Inf-greatest}[OF 0]$ by
auto

thus False using a by simp

qed

ultimately have $f x \in \{a..b\}$ if $x \in \text{space } M$ for x using that by
auto

hence $AE x \text{ in } M. f x \in \{a..b\}$ by simp

thus ?thesis unfolding 1 by (intro Hoeffdings-lemma-bochner assms(1,2,3))
qed

lemma (in prob-space) Hoeffdings-lemma-bochner-3:

assumes expectation $f = 0$

assumes random-variable borel f

assumes $\bigwedge x y. \{x,y\} \subseteq \text{space } M \implies |f x - f y| \leq (c::\text{real})$

shows expectation $(\lambda x. \exp (l * f x)) \leq \exp (l^2 * c^2 / 8)$ (is ?L
 $\leq ?R$)

proof -

consider (a) $l > 0$ | (b) $l = 0$ | (c) $l < 0$

by argo

then show ?thesis

```

proof (cases)
  case a thus ?thesis by (intro Hoeffdings-lemma-bochner-2 assms)
auto
  next
    case b thus ?thesis by simp
  next
    case c
      have ?L = expectation ( $\lambda x. \exp((-l) * (-f x))$ ) by simp
      also have ...  $\leq \exp((-l)^2 * c^2/8)$  using c assms by (intro
Hoeffdings-lemma-bochner-2) auto
      also have ... = ?R by simp
      finally show ?thesis by simp
    qed
  qed

```

Version of *product-sigma-finite.product-integral-singleton* without the condition that $M i$ has to be sigma finite for all i :

```

lemma product-integral-singleton:
  fixes f :: -  $\Rightarrow$  -::{banach, second-countable-topology}
  assumes sigma-finite-measure ( $M i$ )
  assumes  $f \in$  borel-measurable ( $M i$ )
  shows ( $\int x. f (x i) \partial(\text{PiM } \{i\} M)$ ) = ( $\int x. f x \partial(M i)$ ) (is ?L =
?R)
proof -
  define  $M' j =$  (if  $j=i$  then  $M i$  else count-space {undefined})
for j

  interpret product-sigma-finite  $M'$ 
  using assms(1) unfolding product-sigma-finite-def  $M'$ -def
  by (auto intro!:sigma-finite-measure-count-space-finite)

  have ?L =  $\int x. f (x i) \partial(\text{PiM } \{i\} M')$ 
  by (intro Bochner-Integration.integral-cong PiM-cong) (simp-all
add:M'-def)
  also have ... = ( $\int x. f x \partial(M' i)$ )
  using assms(2) by (intro product-integral-singleton) (simp add:M'-def)
  also have ... = ?R
  by (intro Bochner-Integration.integral-cong PiM-cong) (simp-all
add:M'-def)
  finally show ?thesis by simp
qed

```

Version of *product-sigma-finite.product-integral-fold* without the condition that $M i$ has to be sigma finite for all i :

```

lemma product-integral-fold:
  fixes f :: -  $\Rightarrow$  -::{banach, second-countable-topology}
  assumes  $\bigwedge i. i \in I \cup J \implies$  sigma-finite-measure ( $M i$ )
  assumes  $I \cap J = \{\}$ 
  assumes finite  $I$ 

```

assumes *finite J*
assumes *integrable (PiM (I ∪ J) M) f*
shows $(\int x. f x \partial \text{PiM } (I \cup J) M) = (\int x. (\int y. f (\text{merge } I J(x,y)) \partial \text{PiM } J M) \partial \text{PiM } I M)$ (**is** ?L = ?R)
and *integrable (PiM I M) (λx. (∫ y. f (merge I J(x,y)) ∂PiM J M))* (**is** ?I)
and *AE x in PiM I M. integrable (PiM J M) (λy. f (merge I J(x,y)))* (**is** ?T)
proof –
define *M' where M' i = (if i ∈ I ∪ J then M i else count-space {undefined})* **for** *i*

interpret *product-sigma-finite M'*
using *assms(1) unfolding product-sigma-finite-def M'-def*
by (*auto intro!:sigma-finite-measure-count-space-finite*)

interpret *pair-sigma-finite Pi_M I M' Pi_M J M'*
using *assms(3,4) sigma-finite unfolding pair-sigma-finite-def* **by** *blast*

have *0: integrable (Pi_M (I ∪ J) M') f = integrable (Pi_M (I ∪ J) M) f*
by (*intro Bochner-Integration.integrable-cong PiM-cong*) (*simp-all add:M'-def*)

have *?L = (∫ x. f x ∂PiM (I ∪ J) M')*
by (*intro Bochner-Integration.integral-cong PiM-cong*) (*simp-all add:M'-def*)
also have *... = (∫ x. (∫ y. f (merge I J (x,y)) ∂PiM J M') ∂PiM I M')*
using *assms(5) by (intro product-integral-fold assms(2,3,4))* (*simp add:0*)
also have *... = ?R*
by (*intro Bochner-Integration.integral-cong PiM-cong*) (*simp-all add:M'-def*)
finally show *?L = ?R* **by** *simp*

have *integrable (Pi_M (I ∪ J) M') f = integrable (PiM I M' ⊗_M PiM J M') (λx. f (merge I J x))*
using *assms(5) apply (subst distr-merge[OF assms(2,3,4),symmetric])*
by (*intro integrable-distr-eq*) (*simp-all add:0[symmetric]*)
hence *1: integrable (PiM I M' ⊗_M PiM J M') (λx. f (merge I J x))*
using *assms(5) 0* **by** *simp*

hence *integrable (PiM I M') (λx. (∫ y. f (merge I J(x,y)) ∂PiM J M'))* (**is** ?I')
by (*intro integrable-fst'*) *auto*
moreover have *?I' = ?I*

by (intro Bochner-Integration.integrable-cong PiM-cong ext Bochner-Integration.integral-cong)
 (simp-all add:M'-def)
ultimately show ?I
by simp

have AE x in PiM I M'. integrable (PiM J M') (λy. f (merge I J
 (x, y))) (is ?T')
by (intro AE-integrable-fst'[OF 1])
moreover have ?T' = ?T
by (intro arg-cong2[where f=almost-everywhere] PiM-cong ext
 Bochner-Integration.integrable-cong)
 (simp-all add:M'-def)
ultimately show ?T
by simp
qed

lemma product-integral-insert:

fixes f :: - ⇒ -::{banach, second-countable-topology}
assumes $\bigwedge k. k \in \{i\} \cup J \implies \text{sigma-finite-measure } (M k)$
assumes $i \notin J$
assumes finite J
assumes integrable (PiM (insert i J) M) f
shows $(\int x. f x \partial \text{PiM } (\text{insert } i \text{ J}) M) = (\int x. (\int y. f (y(i := x))) \partial \text{PiM } J M) \partial M i$ (is ?L = ?R)
proof –
note meas-cong = iffD1[OF measurable-cong]

have integrable (PiM {i} M) (λx. ($\int y. f (\text{merge } \{i\} J (x,y)) \partial \text{PiM } J M$))
using assms **by** (intro product-integral-fold) auto
hence 0:(λx. ($\int y. f (\text{merge } \{i\} J (x,y)) \partial \text{PiM } J M$)) ∈ borel-measurable
 (PiM {i} M)
using borel-measurable-integrable **by** simp
have 1:(λx. ($\int y. f (y(i := (x i))) \partial \text{PiM } J M$)) ∈ borel-measurable
 (PiM {i} M)
by (intro meas-cong[OF - 0] Bochner-Integration.integral-cong
 arg-cong[where f=f])
 (auto simp add:space-PiM merge-def fun-upd-def PiE-def exten-
 sional-def)
have (λx. ($\int y. f (y(i := (\lambda i \in \{i\}. x) i)) \partial \text{PiM } J M$)) ∈ borel-measurable
 (M i)
by (intro measurable-compose[OF - 1, where f=(λx. (λi ∈ {i}. x))] measurable-restrict) auto
hence 2:(λx. ($\int y. f (y(i := x)) \partial \text{PiM } J M$)) ∈ borel-measurable
 (M i) **by** simp

have ?L = ($\int x. f x \partial \text{PiM } (\{i\} \cup J) M$) **by** simp
also have ... = ($\int x. (\int y. f (\text{merge } \{i\} J (x,y)) \partial \text{PiM } J M) \partial \text{PiM } \{i\} M$)

using *assms(2,4)* **by** (*intro product-integral-fold assms(1,3)*) *auto*
also have ... = $(\int x. (\int y. f (y(i) := (x i))) \partial PiM J M) \partial PiM \{i\} M)$
by (*intro Bochner-Integration.integral-cong refl arg-cong[where f=f]*)
(auto simp add:space-PiM merge-def fun-upd-def PiE-def extensional-def)
also have ... = ?R
using *assms(1,4)* **by** (*intro product-integral-singleton assms(1) 2*)
auto
finally show ?thesis **by** *simp*
qed

lemma *product-integral-insert-rev*:
fixes *f :: - => -::{banach, second-countable-topology}*
assumes $\bigwedge k. k \in \{i\} \cup J \implies \text{sigma-finite-measure } (M k)$
assumes $i \notin J$
assumes *finite J*
assumes *integrable (PiM (insert i J) M) f*
shows $(\int x. f x \partial PiM (insert i J) M) = (\int y. (\int x. f (y(i) := x)) \partial M i) \partial PiM J M$ (**is** ?L = ?R)
proof –
have ?L = $(\int x. f x \partial PiM (J \cup \{i\}) M)$ **by** *simp*
also have ... = $(\int x. (\int y. f (merge J \{i\} (x,y)) \partial PiM \{i\} M) \partial PiM J M)$
using *assms(2,4)* **by** (*intro product-integral-fold assms(1,3)*) *auto*
also have ... = $(\int x. (\int y. f (x(i) := (y i))) \partial PiM \{i\} M) \partial PiM J M)$
unfolding *merge-singleton[OF assms(2)]*
by (*intro Bochner-Integration.integral-cong refl arg-cong[where f=f]*)
(metis PiE-restrict assms(2) restrict-upd space-PiM)
also have ... = ?R
using *assms(1,4)* **by** (*intro Bochner-Integration.integral-cong product-integral-singleton*) *auto*
finally show ?thesis **by** *simp*
qed

lemma *merge-empty[simp]*:
 $merge \{ \} I (y,x) = restrict x I$
 $merge I \{ \} (y,x) = restrict y I$
unfolding *merge-def restrict-def* **by** *auto*

lemma *merge-cong*:
assumes $restrict x1 I = restrict x2 I$
assumes $restrict y1 J = restrict y2 J$
shows $merge I J (x1,y1) = merge I J (x2,y2)$
using *assms* **unfolding** *merge-def restrict-def*
by (*intro ext*) (*smt (verit, best) case-prod-conv*)

lemma *restrict-merge*:

restrict (*merge* $I J x$) $K = \text{merge } (I \cap K) (J \cap K) x$

unfolding *restrict-def merge-def* **by** (*intro ext*) (*auto simp: case-prod-beta*)

lemma *map-prod-measurable*:

assumes $f \in M \rightarrow_M M'$

assumes $g \in N \rightarrow_M N'$

shows *map-prod* $f g \in M \otimes_M N \rightarrow_M M' \otimes_M N'$

using *assms* **by** (*subst measurable-pair-iff*) *simp*

lemma *mc-diarmid-inequality-aux*:

fixes $f :: (\text{nat} \Rightarrow 'a) \Rightarrow \text{real}$

fixes $n :: \text{nat}$

assumes $\bigwedge i. i < n \implies \text{prob-space } (M i)$

assumes $\bigwedge i x y. i < n \implies \{x, y\} \subseteq \text{space } (PiM \{..<n\} M) \implies$

$(\forall j \in \{..<n\} - \{i\}. x j = y j) \implies |f x - f y| \leq c i$

assumes *f-meas*: $f \in \text{borel-measurable } (PiM \{..<n\} M)$ **and** $\varepsilon\text{-gt-0}$:

$\varepsilon > 0$

shows $\mathcal{P}(\omega \text{ in } PiM \{..<n\} M. f \omega - (\int \xi. f \xi \partial PiM \{..<n\} M) \geq$

$\varepsilon) \leq \exp(-2 * \varepsilon^2 / (\sum i < n. (c i)^2))$

(**is** $?L \leq ?R$)

proof –

define h **where** $h k = (\lambda \xi. (\int \omega. f (\text{merge } \{..<k\} \{k..<n\} (\xi, \omega))$

$\partial PiM \{k..<n\} M))$ **for** k

define $t :: \text{real}$ **where** $t = 4 * \varepsilon / (\sum i < n. (c i)^2)$

define V **where** $V i \xi = h (Suc i) \xi - h i \xi$ **for** $i \xi$

obtain $x0$ **where** $x0 : x0 \in \text{space } (PiM \{..<n\} M)$

using *prob-space.not-empty[OF prob-space-PiM]* *assms(1)* **by** *fast-force*

have *delta*: $|f x - f y| \leq c i$ **if** $i < n$

$x \in PiE \{..<n\} (\lambda i. \text{space } (M i))$ $y \in PiE \{..<n\} (\lambda i. \text{space } (M i))$

restrict $x (\{..<n\} - \{i\}) = \text{restrict } y (\{..<n\} - \{i\})$

for $x y i$

proof (*rule assms(2)[OF that(1)]*, *goal-cases*)

case 1

then show *?case* **using** *that(2,3)* **unfolding** *space-PiM* **by** *auto*

next

case 2

then show *?case* **using** *that(4)* **by** (*intro ballI*) (*metis restrict-apply'*)

qed

have *c-ge-0*: $c j \geq 0$ **if** $j < n$ **for** j

proof –

have $0 \leq |f x0 - f x0|$ **by** *simp*
also have $\dots \leq c j$ **using** *x0 unfolding space-PiM* **by** (*intro delta that*) *auto*
finally show *?thesis* **by** *simp*
qed
hence *sum-c-ge-0*: $(\sum_{i < n}. (c i)^2) \geq 0$ **by** (*meson sum-nonneg zero-le-power2*)

hence *t-ge-0*: $t \geq 0$ **using** ε -*gt-0* **unfolding** *t-def* **by** *simp*

note *borel-rules* =
borel-measurable-sum measurable-compose[OF - borel-measurable-exp]
borel-measurable-times

note *int-rules* =
prob-space-PiM assms(1) borel-rules
prob-space.integrable-bounded bounded-intros
have *h-n*: $h n \xi = f \xi$ **if** $\xi \in \text{space } (PiM \{..<n\} M)$ **for** ξ
proof –
have $h n \xi = (\int \omega. f (\lambda i \in \{..<n\}. \xi i) \partial PiM \{ \} M)$
unfolding *h-def* **using** *leD*
by (*intro Bochner-Integration.integral-cong PiM-cong arg-cong*[**where** *f=f*] *restrict-cong*)
auto
also have $\dots = f (\text{restrict } \xi \{..<n\})$
unfolding *PiM-empty* **by** *simp*
also have $\dots = f \xi$
using *that* **unfolding** *space-PiM PiE-def*
by (*simp add: extensional-restrict*)
finally show *?thesis*
by *simp*
qed

have *h-0*: $h 0 \xi = (\int \omega. f \omega \partial PiM \{..<n\} M)$ **for** ξ
unfolding *h-def* **by** (*intro Bochner-Integration.integral-cong PiM-cong refl*)
(simp-all add:space-PiM atLeast0LessThan)

have *h-cong*: $h j \omega = h j \xi$ **if** $\text{restrict } \omega \{..<j\} = \text{restrict } \xi \{..<j\}$
for $j \omega \xi$
using *that* **unfolding** *h-def*
by (*intro Bochner-Integration.integral-cong refl arg-cong*[**where** *f=f*] *merge-cong*) *auto*

have *h-meas*: $h i \in \text{borel-measurable } (PiM I M)$ **if** $i \leq n$ $\{..<i\} \subseteq I$
for $i I$
proof –
have $0: \{..<n\} = \{..<i\} \cup \{i..<n\}$
using *that(1)* **by** *auto*

have 1: $\text{merge } \{..<i\} \{i..<n\} = \text{merge } \{..<i\} \{i..<n\} \circ \text{map-prod}$
 $(\lambda x. \text{restrict } x \{..<i\}) \text{id}$
unfolding $\text{merge-def map-prod-def restrict-def comp-def}$
by $(\text{intro ext}) (\text{auto simp: case-prod-beta}')$

have $\text{merge } \{..<i\} \{i..<n\} \in \text{Pi}_M I M \otimes_M \text{Pi}_M \{i..<n\} M \rightarrow_M$
 $\text{Pi}_M \{..<n\} M$
unfolding 0 **by** $(\text{subst } 1) (\text{intro measurable-comp}[OF - \text{measurable-merge}] \text{map-prod-measurable}$
 $\text{measurable-ident measurable-restrict-subset that}(2))$
hence $(\lambda x. f (\text{merge } \{..<i\} \{i..<n\} x)) \in \text{borel-measurable } (\text{Pi}_M$
 $I M \otimes_M \text{Pi}_M \{i..<n\} M)$
by $(\text{intro measurable-compose}[OF - f-meas])$
thus ?thesis
unfolding $h\text{-def by } (\text{intro sigma-finite-measure.borel-measurable-lebesgue-integral}$
 $\text{prob-space-imp-sigma-finite prob-space-PiM assms}(1)) (\text{auto}$
 $\text{simp: case-prod-beta}')$
qed

have $\text{merge-space-aux: merge } \{..<j\} \{j..<n\} u \in (\Pi_E i \in \{..<n\}. \text{space}$
 $(M i))$
if $j \leq n$ $\text{fst } u \in \text{Pi } \{..<j\} (\lambda i. \text{space } (M i)) \text{snd } u \in \text{Pi } \{j..<n\}$
 $(\lambda i. \text{space } (M i))$
for $u j$
proof –
have $\text{merge } \{..<j\} \{j..<n\} (\text{fst } u, \text{snd } u) \in (\text{PiE } (\{..<j\} \cup \{j..<n\}))$
 $(\lambda i. \text{space } (M i))$
using $\text{that by } (\text{intro iffD2}[OF \text{PiE-cancel-merge}] \text{auto})$
also have $\dots = (\Pi_E i \in \{..<n\}. \text{space } (M i))$
using $\text{that by } (\text{intro arg-cong2}[\text{where } f = \text{PiE}] \text{refl}) \text{auto}$
finally show ?thesis **by** simp
qed

have $\text{merge-space: merge } \{..<j\} \{j..<n\} (u, v) \in (\Pi_E i \in \{..<n\}. \text{space}$
 $(M i))$
if $j \leq n$ $u \in \text{PiE } \{..<j\} (\lambda i. \text{space } (M i)) v \in \text{PiE } \{j..<n\} (\lambda i.$
 $\text{space } (M i))$
for $u v j$
using $\text{that by } (\text{intro merge-space-aux}) (\text{simp-all add: PiE-def})$

have $\text{delta!}: |f x - f y| \leq (\sum i < n. c i)$
if $x \in \text{PiE } \{..<n\} (\lambda i. \text{space } (M i)) y \in \text{PiE } \{..<n\} (\lambda i. \text{space } (M$
 $i))$ **for** $x y$
proof –
define m **where** $m i = \text{merge } \{..<i\} \{i..<n\} (x, y)$ **for** i

have 0: $z \in \text{Pi } I (\lambda i. \text{space } (M i))$ **if** $z \in \text{PiE } \{..<n\} (\lambda i. \text{space}$
 $(M i))$

$I \subseteq \{..<n\}$ for $z I$
using that unfolding *PiE-def by auto*

have \exists : $\{..<Suc\ i\} \cap (\{..<n\} - \{i\}) = \{..<i\}$
 $\{Suc\ i..<n\} \cap (\{..<n\} - \{i\}) = \{Suc\ i..<n\}$
 $\{..<i\} \cap (\{..<n\} - \{i\}) = \{..<i\}$
 $\{i..<n\} \cap (\{..<n\} - \{i\}) = \{Suc\ i..<n\}$
if $i < n$ **for** i
using that by auto

have $|f\ x - f\ y| = |f\ (m\ n) - f\ (m\ 0)|$
using that unfolding *m-def by (simp add:atLeast0LessThan)*
also have $... = |\sum\ i < n. f\ (m\ (Suc\ i)) - f\ (m\ i)|$
by *(subst sum-lessThan-telescope) simp*
also have $... \leq (\sum\ i < n. |f\ (m\ (Suc\ i)) - f\ (m\ i)|)$
by *simp*
also have $... \leq (\sum\ i < n. c\ i)$
using that unfolding *m-def by (intro delta sum-mono merge-space-aux 0 subsetI)*
(simp-all add:restrict-merge 3)
finally show *?thesis*
by *simp*
qed

have $norm\ (f\ x) \leq norm\ (f\ x0) + sum\ c\ \{..<n\}$ **if** $x \in space\ (Pi_M\ \{..<n\}\ M)$ **for** x
proof –
have $|f\ x - f\ x0| \leq sum\ c\ \{..<n\}$
using *x0 that unfolding space-PiM by (intro delta') auto*
thus *?thesis*
by *simp*
qed

hence *f-bounded: bounded (f ' space (PiM {..<n} M))*
by *(intro boundedI[where B=norm (f x0) + (sum i<n. c i)]) auto*

have *f-merge-bounded:*
 $bounded\ ((\lambda\ \omega. (f\ (merge\ \{..<j\}\ \{j..<n\}\ (u, \omega))))\ ' space\ (Pi_M\ \{j..<n\}\ M))$
if $j \leq n$ $u \in PiE\ \{..<j\}$ $(\lambda\ i. space\ (M\ i))$ **for** $u\ j$
proof –
have $(\lambda\ \omega. merge\ \{..<j\}\ \{j..<n\}\ (u, \omega))\ ' space\ (Pi_M\ \{j..<n\}\ M)$
 $\subseteq space\ (Pi_M\ \{..<n\}\ M)$
using that unfolding *space-PiM*
by *(intro image-subsetI merge-space) auto*
thus *?thesis*
by *(subst image-image[of f,symmetric]) (intro bounded-subset[OF f-bounded] image-mono)*
qed

have *f-merge-meas-aux*:
 $(\lambda\omega. f (\text{merge } \{..\<j\} \{j..\<n\} (u, \omega))) \in \text{borel-measurable } (Pi_M \{j..\<n\} M)$
if $j \leq n$ $u \in Pi \{..\<j\} (\lambda i. \text{space } (M i))$ **for** j u
proof –

have $0: \{..\<n\} = \{..\<j\} \cup \{j..\<n\}$
using *that(1)* **by** *auto*

have $1: \text{merge } \{..\<j\} \{j..\<n\} (u, \omega) = \text{merge } \{..\<j\} \{j..\<n\}$
(restrict u {..\<j}, \omega) **for** ω
by *(intro merge-cong) auto*

have $(\lambda\omega. \text{merge } \{..\<j\} \{j..\<n\} (u, \omega)) \in Pi_M \{j..\<n\} M \rightarrow_M Pi_M \{..\<n\} M$
using *that unfolding 0 1*
by *(intro measurable-compose[OF - measurable-merge] measurable-Pair1')*
(simp add:space-PiM)
thus *?thesis*
by *(intro measurable-compose[OF - f-meas])*
qed

have *f-merge-meas*: $(\lambda\omega. f (\text{merge } \{..\<j\} \{j..\<n\} (u, \omega))) \in \text{borel-measurable } (Pi_M \{j..\<n\} M)$
if $j \leq n$ $u \in PiE \{..\<j\} (\lambda i. \text{space } (M i))$ **for** j u
using *that unfolding PiE-def* **by** *(intro f-merge-meas-aux) auto*

have *h-bounded*: $\text{bounded } (h i \text{ 'space } (PiM I M))$
if *h-bounded-assms*: $i \leq n$ $\{..\<i\} \subseteq I$ **for** i I
proof –

have $\text{merge } \{..\<i\} \{i..\<n\} x \in \text{space } (Pi_M \{..\<n\} M)$
if $x \in (\prod_E i \in I. \text{space } (M i)) \times (\prod_E i \in \{i..\<n\}. \text{space } (M i))$ **for** x
using *that h-bounded-assms unfolding space-PiM* **by** *(intro merge-space-aux)*
(auto simp: PiE-def mem-Times-iff)
hence $\text{bounded } ((\lambda x. f (\text{merge } \{..\<i\} \{i..\<n\} x)) \text{ '})$
 $((\prod_E i \in I. \text{space } (M i)) \times (\prod_E i \in \{i..\<n\}. \text{space } (M i)))$
by *(subst image-image[of f, symmetric])*
(intro bounded-subset[OF f-bounded] image-mono image-subsetI)
thus *?thesis*
using *that unfolding h-def*
by *(intro prob-space.finite-measure finite-measure.bounded-int-int-rules)*
(auto simp:space-PiM PiE-def)
qed

have *V-bounded*: $\text{bounded } (V i \text{ 'space } (PiM I M))$

if $i < n$ $\{..<i+1\} \subseteq I$ **for** $i \in I$
using that unfolding $V\text{-def}$ **by** (*intro bounded-intros h-bounded*)
auto

have $V\text{-upd-bounded: bounded } ((\lambda x. V j (\xi(j := x))) \text{ 'space } (M j))$
if $V\text{-upd-bounded-assms: } \xi \in \text{space } (Pi_M \{..<j\} M)$ $j < n$ **for** $j \in \xi$
proof –
have $\xi(j := v) \in \text{space } (Pi_M \{..<j + 1\} M)$ **if** $v \in \text{space } (M j)$
for v
using $V\text{-upd-bounded-assms that unfolding space-PiM PiE-def}$
extensional-def Pi-def by auto
thus *?thesis*
using that unfolding $\text{image-image[of } V j (\lambda x. (\xi(j := x))), \text{symmetric]}$
by (*intro bounded-subset[OF V-bounded[of j {..<j+1}]] that*
image-mono) *auto*
qed

have $h\text{-step: } h j \omega = \int \tau. h (j+1) (\omega (j := \tau)) \partial M j$ (**is** $?L1 =$
 $?R1$)
if $\omega \in \text{space } (Pi_M \{..<j\} M)$ $j < n$ **for** $j \in \omega$
proof –
have $0: (\lambda x. f (\text{merge } \{..<j\} \{j..<n\} (\omega, x))) \in \text{borel-measurable}$
 $(Pi_M \{j..<n\} M)$
using that unfolding $\text{space-PiM by (intro f-merge-meas) auto}$

have $1: \text{insert } j \{Suc j..<n\} = \{j..<n\}$
using that by auto

have $2: \text{bounded } ((\lambda x. (f (\text{merge } \{..<j\} \{j..<n\} (\omega, x)))) \text{ 'space}$
 $(Pi_M \{j..<n\} M))$
using that by (intro f-merge-bounded) (simp-all add: space-PiM)

have $?L1 = (\int \xi. f (\text{merge } \{..<j\} \{j..<n\} (\omega, \xi)) \partial Pi_M (\text{insert } j$
 $\{j+1..<n\} M))$
unfolding $h\text{-def using that by (intro Bochner-Integration.integral-cong}$
 $\text{refl PiM-cong) auto}$
also have $\dots = (\int \tau. (\int \xi. f (\text{merge } \{..<j\} \{j..<n\} (\omega, (\xi(j := \tau))))$
 $\partial Pi_M \{j+1..<n\} M) \partial M j)$
using that(1,2) 0 1 2 by (intro product-integral-insert prob-space-imp-sigma-finite
 assms(1)
 $\text{int-rules f-merge-meas) (simp-all)}$
also have $\dots = ?R1$
using that(2) unfolding $h\text{-def}$
by (intro Bochner-Integration.integral-cong arg-cong[where f=f]
 $\text{ext) (auto simp:merge-def)}$
finally show *?thesis*
by *simp*
qed

have V -meas: $V i \in \text{borel-measurable } (PiM I M)$ **if** $i < n \{..<i+1\}$
 $\subseteq I$ **for** $i I$
unfolding V -def **using that by** (intro borel-measurable-diff h-meas)
auto

have V -upd-meas: $(\lambda x. V j (\xi(j := x))) \in \text{borel-measurable } (M j)$
if $j < n \xi \in \text{space } (PiM \{..<j\} M)$ **for** $j \xi$
using that by (intro measurable-compose[$OF - V$ -meas[**where**
 $I = \text{insert } j \{..<j\}$]])
measurable-component-update) *auto*

have V -cong:
 $V j \omega = V j \xi$ **if** $\text{restrict } \omega \{..<(j+1)\} = \text{restrict } \xi \{..<(j+1)\}$ **for**
 $j \omega \xi$
using that *restrict-subset-eq*[$OF - \text{that}$] **unfolding** V -def
by (intro *arg-cong2*[**where** $f = (-)$] *h-cong*) *simp-all*

have $\text{exp-}V$: $(\int \omega. V j (\xi(j := \omega)) \partial M j) = 0$ (**is** $?L1 = 0$)
if $j < n \xi \in \text{space } (PiM \{..<j\} M)$ **for** $j \xi$
proof –

have $\text{fun-upd } \xi j$ ‘ $\text{space } (M j) \subseteq \text{space } (PiM (\text{insert } j \{..<j\}) M)$
using that **unfolding** *space-PiM* **by** (intro *image-subsetI* *PiE-fun-upd*)
auto
hence 0 :bounded $((\lambda x. h (Suc j) (\xi(j := x)))$ ‘ $\text{space } (M j)$)
unfolding *image-image*[of $h (Suc j) \lambda x. \xi(j := x)$, *symmetric*]
using that
by (intro *bounded-subset*[OF *h-bounded*[**where** $i = j + 1$ **and** $I = \{..<j+1\}$]])
image-mono)
(auto simp:lessThan-Suc)

have 1 : $(\lambda x. h (Suc j) (\xi(j := x))) \in \text{borel-measurable } (M j)$
using *h-meas* **that by** (intro *measurable-compose*[$OF - h$ -meas[**where**
 $I = \text{insert } j \{..<j\}$]])
measurable-component-update) *auto*

have $?L1 = (\int \omega. h (Suc j) (\xi(j := \omega)) - h j \xi \partial M j)$
unfolding V -def
by (intro *Bochner-Integration.integral-cong* *arg-cong2*[**where**
 $f = (-)$] *refl h-cong*) *auto*
also have $\dots = (\int \omega. h (Suc j) (\xi(j := \omega)) \partial M j) - (\int \omega. h j \xi \partial M j)$
using that by (intro *Bochner-Integration.integral-diff* *int-rules 0 1*) *auto*
also have $\dots = 0$
using that (1) *assms(1) prob-space.prob-space* **unfolding** *h-step*[OF
 $\text{that}(2,1)$] **by** *auto*
finally show $?thesis$
by *simp*

qed

have $var-V: |V j x - V j y| \leq c j$ (**is** $?L1 \leq ?R1$)
if $var-V-assms: j < n \{x,y\} \subseteq space (PiM \{..<j+1\} M)$
 $restrict x \{..<j\} = restrict y \{..<j\}$ **for** $x y j$

proof –

have $x-ran: x \in PiE \{..<j+1\} (\lambda i. space (M i))$ **and** $y-ran: y \in PiE \{..<j+1\} (\lambda i. space (M i))$
using $that(2)$ **by** $(simp-all add:space-PiM)$

have $0: j+1 \leq n$
using $that$ **by** $simp$

have $?L1 = |h (Suc j) x - h j y - (h (Suc j) y - h j y)|$
unfolding $V-def$ **by** $(intro arg-cong[where f=abs] arg-cong2[where f=(-)] refl h-cong that)$

also have $... = |h (j+1) x - h (j+1) y|$
by $simp$

also have $... =$

$|(\int \omega. f(merge \{..<j+1\} \{j+1..<n\} (x,\omega)) - f(merge \{..<j+1\} \{j+1..<n\} (y,\omega))) \partial PiM \{j+1..<n\} M|$

using $that$ **unfolding** $h-def$ **by** $(intro arg-cong[where f=abs] f-merge-meas[OF 0] x-ran$

$Bochner-Integration.integral-diff[symmetric] int-rules f-merge-bounded[OF 0] y-ran) auto$

also have $... \leq$

$(\int \omega. |f(merge \{..<j+1\} \{j+1..<n\} (x,\omega)) - f(merge \{..<j+1\} \{j+1..<n\} (y,\omega))| \partial PiM \{j+1..<n\} M)$

by $(intro integral-abs-bound)$

also have $... \leq (\int \omega. c j \partial PiM \{j+1..<n\} M)$

proof $(intro Bochner-Integration.integral-mono' delta int-rules c-ge-0 ballI merge-space 0)$

fix ω **assume** $\omega \in space (PiM \{j+1..<n\} M)$

have $\{..<j+1\} \cap (\{..<n\} - \{j\}) = \{..<j\}$

using $that$ **by** $auto$

thus $restrict (merge \{..<j+1\} \{j+1..<n\} (x, \omega)) (\{..<n\} - \{j\})$

$=$

$restrict (merge \{..<j+1\} \{j+1..<n\} (y, \omega)) (\{..<n\} - \{j\})$

using $that(1,3) less-antisym$ **unfolding** $restrict-merge$ **by** $(intro merge-cong refl) auto$

qed $(simp-all add: space-PiM that(1) x-ran[simplified] y-ran[simplified])$

also have $... = c j$

by $(auto intro!: prob-space. prob-space-PiM assms(1))$

finally show $?thesis$ **by** $simp$

qed

have $f \xi - (\int \omega. f \omega \partial (PiM \{..<n\} M)) = (\sum i < n. V i \xi)$ **if** $\xi \in space (PiM \{..<n\} M)$ **for** ξ

using $that$ **unfolding** $V-def$ **by** $(subst sum-lessThan-telescope)$

(simp add: h-0 h-n)
hence $?L = \mathcal{P}(\xi \text{ in } PiM \{..<n\} M. (\sum i < n. V i \xi) \geq \varepsilon)$
by (*intro arg-cong2[where f=measure] refl Collect-restr-cong arg-cong2[where f=(≤)] auto*)
also have $\dots \leq \mathcal{P}(\xi \text{ in } PiM \{..<n\} M. \exp(t * (\sum i < n. V i \xi)) \geq \exp(t * \varepsilon))$
proof (*intro finite-measure.finite-measure-mono subsetI prob-space.finite-measure int-rules*)
show $\{\xi \in \text{space } (PiM \{..<n\} M). \exp(t * \varepsilon) \leq \exp(t * (\sum i < n. V i \xi))\} \in \text{sets } (PiM \{..<n\} M)$
using *V-meas by measurable*
qed (*auto intro!:mult-left-mono[OF - t-ge-0]*)
also have $\dots \leq (\int \xi. \exp(t * (\sum i < n. V i \xi)) \partial PiM \{..<n\} M) / \exp(t * \varepsilon)$
by (*intro integral-Markov-inequality-measure[where A={}] int-rules V-bounded V-meas*) *auto*
also have $\dots = \exp(t^2 * (\sum i \in \{n..<n\}. c i^2) / 8 - t * \varepsilon) * (\int \xi. \exp(t * (\sum i < n. V i \xi)) \partial PiM \{..<n\} M)$
by (*simp add:exp-minus inverse-eq-divide*)
also have $\dots \leq \exp(t^2 * (\sum i \in \{0..<n\}. c i^2) / 8 - t * \varepsilon) * (\int \xi. \exp(t * (\sum i < 0. V i \xi)) \partial PiM \{..<0\} M)$
proof (*rule ineq-chain*)
fix *j* **assume** *a:j < n*
let $?L1 = \exp(t^2 * (\sum i=j+1..<n. (c i)^2) / 8 - t * \varepsilon)$
let $?L2 = ?L1 * (\int \xi. \exp(t * (\sum i < j+1. V i \xi)) \partial PiM \{..<j+1\} M)$
note $V\text{-upd-meas} = V\text{-upd-meas}[OF a]$
have $?L2 = ?L1 * (\int \xi. \exp(t * (\sum i < j. V i \xi)) * \exp(t * V j \xi) \partial PiM (\text{insert } j \{..<j\}) M)$
by (*simp add:algebra-simps exp-add lessThan-Suc*)
also have $\dots = ?L1 * (\int \xi. (\int \omega. \exp(t * (\sum i < j. V i (\xi(j := \omega)))) * \exp(t * V j (\xi(j := \omega))) \partial M j) \partial PiM \{..<j\} M)$
using *a* **by** (*intro product-integral-insert-rev arg-cong2[where f=(*)] int-rules prob-space-imp-sigma-finite V-bounded V-meas*) *auto*
also have $\dots = ?L1 * (\int \xi. (\int \omega. \exp(t * (\sum i < j. V i \xi)) * \exp(t * V j (\xi(j := \omega))) \partial M j) \partial PiM \{..<j\} M)$
by (*intro arg-cong2[where f=(*)] Bochner-Integration.integral-cong arg-cong[where f=exp] sum.cong V-cong restrict-fupd*) *auto*
also have $\dots = ?L1 * (\int \xi. \exp(t * (\sum i < j. V i \xi)) * (\int \omega. \exp(t * V j (\xi(j := \omega))) \partial M j) \partial PiM \{..<j\} M)$
using *a* **by** (*intro arg-cong2[where f=(*)] Bochner-Integration.integral-cong refl Bochner-Integration.integral-mult-right V-upd-meas V-upd-bounded int-rules*) *auto*
also have $\dots \leq ?L1 * \int \xi. \exp(t * (\sum i < j. V i \xi)) * \exp(t^2 * c$

$j^2/8) \partial PiM \{..<j\} M$
proof (intro mult-left-mono integral-mono)
fix ξ **assume** $c:\xi \in \text{space } (PiM \{..<j\} M)$
hence $b:\xi \in PiE \{..<j\} (\lambda i. \text{space } (M i))$
unfolding space-PiM **by** simp
moreover **have** $\xi(j := v) \in PiE \{..<j+1\} (\lambda i. \text{space } (M i))$ **if**
 $v \in \text{space } (M j)$ **for** v
using b **that** **unfolding** PiE-def extensional-def Pi-def **by** auto
ultimately **show** $LINT \omega | M j. \exp(t * V j (\xi(j := \omega))) \leq \exp$
 $(t^2 * (c j)^2 / 8)$
using V-upd-meas[OF c]
by (intro prob-space.Hoeffdings-lemma-bochner-3 exp-V var-V a
int-rules)
(auto simp: space-PiM)
next
show integrable (PiM {..<j} M) $(\lambda x. \exp(t * (\sum i < j. V i x))) * \exp$
 $(t^2 * (c j)^2 / 8)$
using a **by** (intro int-rules V-bounded V-meas) auto
qed auto
also **have** ... = ?L1 * (($\int \xi. \exp(t * (\sum i < j. V i \xi)) \partial PiM \{..<j\}$
 $M) * \exp(t^2 * c j^2 / 8)$)
proof (subst Bochner-Integration.integral-mult-left)
show integrable (PiM {..<j} M) $(\lambda \xi. \exp(t * (\sum i < j. V i \xi)))$
using a **by** (intro int-rules V-bounded V-meas) auto
qed auto
also **have** ... =
 $\exp(t^2 * (\sum i \in \text{insert } j \{j+1..<n\}. (c i)^2) / 8 - t * \varepsilon) * (\int \xi. \exp(t * (\sum i < j. V i \xi)) \partial PiM \{..<j\} M)$
by (simp-all add:exp-add[symmetric] field-simps)
also **have** ... = $\exp(t^2 * (\sum i = j..<n. (c i)^2) / 8 - t * \varepsilon) * (\int \xi. \exp(t * (\sum i < j. V i \xi)) \partial PiM \{..<j\} M)$
using a **by** (intro arg-cong2[where f=(*)] arg-cong[where
 $f = \exp$] refl arg-cong2
[where f=(-)] arg-cong2[where f=(/)] sum.cong) auto
finally **show** ?L2 $\leq \exp(t^2 * (\sum i = j..<n. (c i)^2) / 8 - t * \varepsilon) * (\int \xi. \exp$
 $(t * (\sum i < j. V i \xi)) \partial PiM \{..<j\} M)$
by simp
qed
also **have** ... = $\exp(t^2 * (\sum i < n. c i^2) / 8 - t * \varepsilon)$ **by** (simp add:PiM-empty
atLeast0LessThan)
also **have** ... = $\exp(t * ((t * (\sum i < n. c i^2) / 8) - \varepsilon))$ **by** (simp
add:algebra-simps power2-eq-square)
also **have** ... = $\exp(t * (-\varepsilon / 2))$ **using** sum-c-ge-0 **by** (auto simp
add:divide-simps t-def)
also **have** ... = ?R **unfolding** t-def **by** (simp add:field-simps power2-eq-square)
finally **show** ?thesis **by** simp
qed

theorem mc-diarmid-inequality-distr:

fixes $f :: ('i \Rightarrow 'a) \Rightarrow \text{real}$
assumes $\text{finite } I$
assumes $\bigwedge i. i \in I \Longrightarrow \text{prob-space } (M\ i)$
assumes $\bigwedge i\ x\ y. i \in I \Longrightarrow \{x, y\} \subseteq \text{space } (PiM\ I\ M) \Longrightarrow (\forall j \in I - \{i\}. x\ j = y\ j) \Longrightarrow |f\ x - f\ y| \leq c\ i$
assumes $f\text{-meas}: f \in \text{borel-measurable } (PiM\ I\ M)$ **and** $\varepsilon\text{-gt-0}: \varepsilon > 0$
shows $\mathcal{P}(\omega \text{ in } PiM\ I\ M. f\ \omega - (\int \xi. f\ \xi\ \partial PiM\ I\ M) \geq \varepsilon) \leq \exp(-(\mathcal{L} * \varepsilon^{\wedge 2}) / (\sum i \in I. (c\ i)^{\wedge 2}))$
(is ?L ≤ ?R)
proof –
define n **where** $n = \text{card } I$
let $?q = \text{from-nat-into } I$
let $?r = \text{to-nat-on } I$
let $?f = (\lambda \xi. f\ (\lambda i \in I. \xi\ (?r\ i)))$

have $q: \text{bij-betw } ?q\ \{\dots < n\}\ I$ **unfolding** $n\text{-def}$ **by** $(\text{intro } \text{bij-betw-from-nat-into-finite } \text{assms}(1))$
have $r: \text{bij-betw } ?r\ I\ \{\dots < n\}$ **unfolding** $n\text{-def}$ **by** $(\text{intro } \text{to-nat-on-finite } \text{assms}(1))$

have $[\text{simp}]: ?q\ (?r\ x) = x$ **if** $x \in I$ **for** x
by $(\text{intro } \text{from-nat-into-to-nat-on that countable-finite } \text{assms}(1))$

have $[\text{simp}]: ?r\ (?q\ x) = x$ **if** $x < n$ **for** x
using $\text{bij-betw-imp-surj-on}[OF\ r]$ **that by** $(\text{intro } \text{to-nat-on-from-nat-into } \text{auto})$

have $a: \bigwedge i. i \in \{\dots < n\} \Longrightarrow \text{prob-space } ((M \circ ?q)\ i)$
unfolding comp-def **by** $(\text{intro } \text{assms}(2)\ \text{bij-betw-apply}[OF\ q])$

have $b: PiM\ I\ M = PiM\ I\ (\lambda i. (M \circ ?q)\ (?r\ i))$ **by** $(\text{intro } PiM\text{-cong})$
 $(\text{simp-all } \text{add: comp-def})$
also have $\dots = \text{distr } (PiM\ \{\dots < n\}\ (M \circ ?q))\ (PiM\ I\ (\lambda i. (M \circ ?q)\ (?r\ i)))\ (\lambda \omega. \lambda n \in I. \omega\ (?r\ n))$
using r **unfolding** bij-betw-def **by** $(\text{intro } \text{distr-PiM-reindex}[\text{symmetric}])$
 $a)$ auto
finally have $c: PiM\ I\ M = \text{distr } (PiM\ \{\dots < n\}\ (M \circ ?q))\ (PiM\ I\ (\lambda i. (M \circ ?q)\ (?r\ i)))\ (\lambda \omega. \lambda n \in I. \omega\ (?r\ n))$
by simp

have $d: (\lambda n \in I. x\ (?r\ n)) \in \text{space } (PiM\ I\ M)$ **if** $\lambda x \in \text{space } (PiM\ \{\dots < n\}\ (M \circ ?q))$ **for** x
proof –
have $x\ (?r\ i) \in \text{space } (M\ i)$ **if** $i \in I$ **for** i
proof –
have $?r\ i \in \{\dots < n\}$ **using** $\text{bij-betw-apply}[OF\ r]$ **that by** simp
hence $x\ (?r\ i) \in \text{space } ((M \circ ?q)\ (?r\ i))$ **using** $\text{that } \lambda PiE\text{-mem}$
unfolding space-PiM **by** blast
thus $?thesis$ **using** $\text{that unfolding comp-def by simp}$

```

qed
thus ?thesis unfolding space-PiM PiE-def by auto
qed

have ?L =  $\mathcal{P}(\omega \text{ in } PiM \{..<n\} (M \circ ?q). ?f \omega - (\int \xi. f \xi \partial PiM I M) \geq \varepsilon)$ 
proof (subst c, subst measure-distr, goal-cases)
  case 1 thus ?case
    by (intro measurable-restrict measurable-component-singleton
    bij-betw-apply[OF r])
  next
    case 2 thus ?case unfolding b[symmetric] by (intro measurable-sets-Collect[OF f-meas]) auto
  next
    case 3 thus ?case using d by (intro arg-cong2[where f=measure] refl) (auto simp:vimage-def)
qed
also have ... =  $\mathcal{P}(\omega \text{ in } PiM \{..<n\} (M \circ ?q). ?f \omega - (\int \xi. ?f \xi \partial PiM \{..<n\} (M \circ ?q)) \geq \varepsilon)$ 
proof (subst c, subst integral-distr, goal-cases)
  case (1  $\omega$ ) thus ?case
    by (intro measurable-restrict measurable-component-singleton
    bij-betw-apply[OF r])
  next
    case (2  $\omega$ ) thus ?case unfolding b[symmetric] by (rule f-meas)
  next
    case 3 thus ?case by simp
qed
also have ...  $\leq \exp(-(\mathcal{L} * \varepsilon^{\wedge 2}) / (\sum i < n. (c (?q i))^{\wedge 2}))$ 
proof (intro mc-diarmid-inequality-aux  $\varepsilon$ -gt-0, goal-cases)
  case (1  $i$ ) thus ?case by (intro a) auto
next
  case (2  $i$   $x$   $y$ )
  have  $x (?r j) = y (?r j)$  if  $j \in I - \{?q i\}$  for  $j$ 
  proof -
    have  $?r j \in \{..<n\} - \{i\}$  using that bij-betw-apply[OF r] by auto
    thus ?thesis using 2 by simp
  qed
  hence  $\forall j \in I - \{?q i\}. (\lambda i \in I. x (?r i)) j = (\lambda i \in I. y (?r i)) j$  by auto
  thus ?case using 2 d by (intro assms(3) bij-betw-apply[OF q]) auto
next
  case 3
  have  $(\lambda x. x (?r i)) \in PiM \{..<n\} (M \circ ?q) \rightarrow_M M i$  if  $i \in I$  for  $i$ 
  proof -
    have  $0 : M i = (M \circ ?q) (?r i)$  using that by (simp add: comp-def)
    show ?thesis unfolding 0 by (intro measurable-component-singleton

```

bij-betw-apply[*OF r*] *that*)
qed
thus *?case by* (*intro measurable-compose*[*OF - f-meas*] *measurable-restrict*)
qed
also have ... = *?R by* (*subst sum.reindex-bij-betw*[*OF q*]) *simp*
finally show *?thesis by simp*
qed

lemma (*in prob-space*) *mc-diarmid-inequality-classic*:

fixes *f :: ('i ⇒ 'a) ⇒ real*
assumes *finite I*
assumes *indep-vars N X I*
assumes $\bigwedge i x y. i \in I \implies \{x, y\} \subseteq \text{space } (PiM I N) \implies (\forall j \in I - \{i\}. x j = y j) \implies |f x - f y| \leq c i$
assumes *f-meas: f ∈ borel-measurable (PiM I N) and ε-gt-0: ε > 0*
shows $\mathcal{P}(\omega \text{ in } M. f (\lambda i \in I. X i \omega) - (\int \xi. f (\lambda i \in I. X i \xi) \partial M) \geq \varepsilon) \leq \exp(-2 * \varepsilon^2 / (\sum i \in I. (c i)^2))$
(is ?L ≤ ?R)

proof –

note *indep-imp = iffD1*[*OF indep-vars-iff-distr-eq-PiM'*]
let *?O = λi. distr M (N i) (X i)*
have *a: distr M (PiM I N) (λx. λi ∈ I. X i x) = PiM I ?O*
using *assms(2) unfolding indep-vars-def by (intro indep-imp*[*OF - assms(2)*]) *auto*

have *b: space (PiM I ?O) = space (PiM I N)*
by (*metis (no-types, lifting) a space-distr*)

have $(\lambda i \in I. X i \omega) \in \text{space } (PiM I N)$ **if** $\omega \in \text{space } M$ **for** ω
using *assms(2) that unfolding indep-vars-def measurable-def space-PiM by auto*

hence $?L = \mathcal{P}(\omega \text{ in } M. (\lambda i \in I. X i \omega) \in \text{space } (PiM I N) \wedge f (\lambda i \in I. X i \omega) - (\int \xi. f (\lambda i \in I. X i \xi) \partial M) \geq \varepsilon)$

by (*intro arg-cong2*[*where f=measure*] *Collect-restr-cong refl*) *auto*
also have ... = $\mathcal{P}(\omega \text{ in } \text{distr } M (PiM I N) (\lambda x. \lambda i \in I. X i x). f \omega - (\int \xi. f (\lambda i \in I. X i \xi) \partial M) \geq \varepsilon)$

proof (*subst measure-distr, goal-cases*)

case 1 thus *?case using assms(2) unfolding indep-vars-def by (intro measurable-restrict) auto*

next

case 2 thus *?case unfolding space-distr by (intro measurable-sets-Collect*[*OF f-meas*]) *auto*

next

case 3 thus *?case by (simp-all add: Int-def conj-commute)*

qed

also have ... = $\mathcal{P}(\omega \text{ in } PiM I ?O. f \omega - (\int \xi. f (\lambda i \in I. X i \xi) \partial M) \geq \varepsilon)$

```

    unfolding a by simp
  also have ... =  $\mathcal{P}(\omega \text{ in } PiM I ?O. f \omega - (\int \xi. f \xi \partial \text{ distr } M (PiM I N) (\lambda x. \lambda i \in I. X i x)) \geq \varepsilon)$ 
  proof (subst integral-distr[OF - f-meas], goal-cases)
    case (1  $\omega$ ) thus ?case using assms(2) unfolding indep-vars-def
  by (intro measurable-restrict) auto
  next
    case 2 thus ?case by simp
  qed
  also have ... =  $\mathcal{P}(\omega \text{ in } PiM I ?O. f \omega - (\int \xi. f \xi \partial PiM I ?O) \geq \varepsilon)$ 
  unfolding a by simp
  also have ...  $\leq ?R$ 
    using f-meas assms(2) b unfolding indep-vars-def
    by (intro mc-diarmid-inequality-distr prob-space-distr assms(1)
 $\varepsilon$ -gt-0 assms(3)) auto
  finally show ?thesis by simp
qed

end

```

7 Paley-Zygmund Inequality

This section proves slight improvements of the Paley-Zygmund Inequality [7]. Unfortunately, the improvements are on Wikipedia with no citation.

```
theory Paley-Zygmund-Inequality
```

```
imports Lp.Lp
```

```
begin
```

```
context prob-space
```

```
begin
```

```
theorem paley-zygmund-inequality-holder:
```

```
assumes p:  $1 < (p::real)$ 
```

```
assumes rv: random-variable borel Z
```

```
assumes intZp: integrable M ( $\lambda z. |Z z| \text{ powr } p$ )
```

```
assumes t:  $\vartheta \leq 1$ 
```

```
assumes ZAEpos:  $\forall z \text{ in } M. Z z \geq 0$ 
```

```
shows
```

```
( $\text{expectation } (\lambda x. |Z x - \vartheta * \text{expectation } Z| \text{ powr } p) \text{ powr } (1 / (p-1))) *$ 
```

```
 $\text{prob } \{z \in \text{space } M. Z z > \vartheta * \text{expectation } Z\}$ 
```

```
 $\geq ((1-\vartheta) \text{ powr } (p / (p-1))) * \text{expectation } Z \text{ powr } (p / (p-1))$ 
```

```
proof -
```

```
have intZ: integrable M Z
```

```
apply (subst bound-L1-Lp[OF - rv intZp])
```

```
using p by auto
```

```

define eZ where eZ = expectation Z
have eZ ≥ 0
  unfolding eZ-def
  using ZAEpos intZ integral-ge-const prob-Collect-eq-1 by auto

have ezp: expectation (λx. |Z x - ∅ * eZ| powr p) ≥ 0
  by (meson Bochner-Integration.integral-nonneg powr-ge-zero)

have expectation (λz. Z z - ∅ * eZ) = expectation (λz. Z z + (- ∅
* eZ))
  by auto
moreover have ... = expectation Z + expectation (λz. - ∅ * eZ)
  apply (subst Bochner-Integration.integral-add)
  using intZ by auto
moreover have ... = eZ + (- ∅ * eZ)
  apply (subst lebesgue-integral-const)
  using eZ-def prob-space by auto
ultimately have *: expectation (λz. Z z - ∅ * eZ) = eZ - ∅ * eZ
  by linarith

have ev: {z ∈ space M. ∅ * eZ < Z z} ∈ events
  using rv unfolding borel-measurable-iff-greater
  by auto

define q where q = p / (p-1)

have sqI:(indicat-real E x) powr q = indicat-real E (x::'a) for E x
  unfolding q-def
  by (metis indicator-simps(1) indicator-simps(2) powr-0 powr-one-eq-one)

have bm1: (λz. (Z z - ∅ * eZ)) ∈ borel-measurable M
  using borel-measurable-const borel-measurable-diff rv by blast
have bm2: (λz. indicat-real {z ∈ space M. Z z > ∅ * eZ} z) ∈
borel-measurable M
  using borel-measurable-indicator ev by blast
have integrable M (λx. |Z x + (-∅ * eZ)| powr p)
  apply (intro Minkowski-inequality[OF - rv - intZp])
  using p by auto
then have int1: integrable M (λx. |Z x - ∅ * eZ| powr p)
  by auto

have integrable M
(λx. 1 * indicat-real {z ∈ space M. ∅ * eZ < Z z} x)
  apply (intro integrable-real-mult-indicator[OF ev])
  by auto

then have int2: integrable M
(λx. |indicat-real {z ∈ space M. ∅ * eZ < Z z} x| powr q)
  by (auto simp add: sqI )

```

```

have pq:p > (0::real) q > 0 1/p + 1/q = 1
  unfolding q-def using p by (auto simp:divide-simps)
from Holder-inequality[OF pq bm1 bm2 int1 int2]
have hi: expectation (λx. (Z x - v * eZ) * indicat-real {z ∈ space
M. v * eZ < Z z} x)
  ≤ expectation (λx. |Z x - v * eZ| powr p) powr (1 / p) *
  expectation (λx. |indicat-real {z ∈ space M. v * eZ < Z z} x|
powr q) powr (1 / q)
  by auto

have eZ - v * eZ ≤
  expectation (λz. (Z z - v * eZ) * indicat-real {z ∈ space M. Z z
> v * eZ} z)
  unfolding *[symmetric]
  apply (intro integral-mono)
  using intZ ev apply auto[1]
  apply (auto intro!: integrable-real-mult-indicator simp add: intZ
ev)[1]
  unfolding indicator-def of-bool-def
  by (auto simp add: mult-nonneg-nonpos2)

also have ... ≤
  expectation (λx. |Z x - v * eZ| powr p) powr (1 / p) *
  expectation (λx. indicat-real {z ∈ space M. v * eZ < Z z} x)
powr (1 / q)
  using hi by (auto simp add: sqI)

finally have eZ - v * eZ ≤
  expectation (λx. |Z x - v * eZ| powr p) powr (1 / p) *
  expectation (λx. indicat-real {z ∈ space M. v * eZ < Z z} x) powr
(1 / q)
  by auto

then have (eZ - v * eZ) powr q ≤
  (expectation (λx. |Z x - v * eZ| powr p) powr (1 / p) *
  expectation (λx. indicat-real {z ∈ space M. v * eZ < Z z} x) powr
(1 / q)) powr q
  by (smt (verit, ccfv-SIG) ⟨0 ≤ eZ⟩ mult-left-le-one-le powr-mono2
pq(2) right-diff-distrib' t zero-le-mult-iff)

also have ... =
  (expectation (λx. |Z x - v * eZ| powr p) powr (1 / p)) powr q *
  (expectation (λx. indicat-real {z ∈ space M. v * eZ < Z z} x)
powr (1 / q)) powr q
  using powr-ge-zero powr-mult by presburger
also have ... =
  (expectation (λx. |Z x - v * eZ| powr p) powr (1 / p)) powr q *
  (expectation (λx. indicat-real {z ∈ space M. v * eZ < Z z} x))

```


by (*smt* (*verit*, *ccfv-SIG*) *Bochner-Integration.integral-nonneg divide-le-eq-1-pos indicator-pos-le nonzero-eq-divide-eq p powr-one powr-powr q-def*)
also have ... =
 (*expectation* ($\lambda x. |Z x - \vartheta * eZ| \text{ powr } p$) *powr* ($1 / (p-1)$)) *
 (*expectation* ($\lambda x. \text{indicat-real } \{z \in \text{space } M. \vartheta * eZ < Z z\} x$))
by (*smt* (*verit*, *ccfv-threshold*) *divide-divide-eq-right divide-self-if p powr-powr q-def times-divide-eq-left*)
also have ... =
 (*expectation* ($\lambda x. |Z x - \vartheta * eZ| \text{ powr } p$) *powr* ($1 / (p-1)$)) *
prob $\{z \in \text{space } M. Z z > \vartheta * eZ\}$
by (*simp add: ev*)

finally have 1: ($eZ - \vartheta * eZ$) *powr* $q \leq$
 (*expectation* ($\lambda x. |Z x - \vartheta * eZ| \text{ powr } p$) *powr* ($1 / (p-1)$)) *
prob $\{z \in \text{space } M. Z z > \vartheta * eZ\}$ **by** *linarith*

have ($eZ - \vartheta * eZ$) *powr* $q = ((1 - \vartheta) * eZ)$ *powr* q
by (*simp add: mult.commute right-diff-distrib*)
also have ... = $(1 - \vartheta)$ *powr* $q * eZ$ *powr* q
by (*simp add: <0 ≤ eZ> powr-mult t*)
finally show ?thesis **using** 1 *eZ-def q-def* **by force**
qed

corollary *paley-zygmund-inequality:*

assumes *rv: random-variable borel Z*

assumes *intZsq: integrable M* ($\lambda z. (Z z)^2$)

assumes *t: $\vartheta \leq 1$*

assumes *Zpos: $\bigwedge z. z \in \text{space } M \implies Z z \geq 0$*

shows

(*variance* $Z + (1-\vartheta)^2 * (\text{expectation } Z)^2$) *
prob $\{z \in \text{space } M. Z z > \vartheta * \text{expectation } Z\}$
 $\geq (1-\vartheta)^2 * (\text{expectation } Z)^2$

proof –

have *ZAEpos: AE z in M. Z z ≥ 0*

by (*simp add: Zpos*)

define *p where* $p = (2::\text{real})$

have *p1: 1 < p using p-def* **by auto**

have *integrable M* ($\lambda z. |Z z| \text{ powr } p$) **unfolding** *p-def*
using *intZsq* **by auto**

from *paley-zygmund-inequality-holder[OF p1 rv this t ZAEpos]*

have $(1 - \vartheta) \text{ powr } (p / (p - 1)) * (\text{expectation } Z \text{ powr } (p / (p - 1)))$
 $\leq \text{expectation } (\lambda x. |Z x - \vartheta * \text{expectation } Z| \text{ powr } p) \text{ powr } (1 / (p - 1)) *$
prob $\{z \in \text{space } M. \vartheta * \text{expectation } Z < Z z\}$.

```

then have hi:  $(1 - \vartheta)^2 * (\text{expectation } Z)^2$ 
   $\leq \text{expectation } (\lambda x. (Z x - \vartheta * \text{expectation } Z)^2) *$ 
   $\text{prob } \{z \in \text{space } M. \vartheta * \text{expectation } Z < Z z\}$ 
unfolding p-def by (auto simp add: Zpos t)

have intZ: integrable M Z
  apply (subst square-integrable-imp-integrable[OF rv intZsq])
  by auto

define eZ where eZ = expectation Z
have eZ  $\geq 0$ 
  unfolding eZ-def
  using Bochner-Integration.integral-nonneg Zpos by blast

have exp: expectation  $(\lambda x. |Z x - \vartheta * eZ| \text{ powr } p) \geq 0$ 
  by (meson Bochner-Integration.integral-nonneg powr-ge-zero)

have expectation  $(\lambda z. Z z - \vartheta * eZ) = \text{expectation } (\lambda z. Z z + (- \vartheta$ 
* eZ))
  by auto
also have ... = expectation Z + expectation  $(\lambda z. - \vartheta * eZ)$ 
  apply (subst Bochner-Integration.integral-add)
  using intZ by auto
also have ... = eZ +  $(- \vartheta * eZ)$ 
  apply (subst lebesgue-integral-const)
  using eZ-def prob-space by auto
finally have *: expectation  $(\lambda z. Z z - \vartheta * eZ) = eZ - \vartheta * eZ$ 
  by linarith
have variance Z =
  variance  $(\lambda z. (Z z - \vartheta * eZ))$ 
  using * eZ-def by auto
also have ... =
  expectation  $(\lambda z. (Z z - \vartheta * eZ)^2)$ 
   $- (\text{expectation } (\lambda x. Z x - \vartheta * eZ))^2$ 
  apply (subst variance-eq)
  by (auto simp add: intZ power2-diff intZsq)
also have ... = expectation  $(\lambda z. (Z z - \vartheta * eZ)^2) - ((1-\vartheta)^2 * eZ^2)$ 
  unfolding * by (auto simp: algebra-simps power2-eq-square)
finally have veq: expectation  $(\lambda z. (Z z - \vartheta * eZ)^2) = (\text{variance } Z$ 
+  $(1-\vartheta)^2 * eZ^2)$ 
  by linarith
thus ?thesis
  using hi by (simp add: eZ-def)
qed

end

end

```

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