

Formalisation and Analysis of Component Dependencies

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Abstract

This set of theories presents a formalisation in Isabelle/HOL [3] of data dependencies between components. The approach allows to analyse system structure oriented towards efficient checking of system: it aims at elaborating for a concrete system, which parts of the system (or system model) are necessary to check a given property.

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1 Introduction

In general, we don't need complete information about the system as to check its certain property. An additional information about the system can slow the whole process down or even make it infeasible. In this theory we define constraints that allow to find/check the minimal model (and the minimal extent of the system) needed to verify a specific property. Our approach focuses on data dependencies between system components. Dependencies' analysis results in a decomposition that gives rise to a logical system architecture, which is the most appropriate for the case of remote monitoring, testing and/or verification.

Let $CSet$ be a set of components on a certain abstraction level L of logical architecture (i.e. level of refinement/decomposition, data type *AbstrLevelsID* in our Isabelle formalisation). We denote the sets of input and output streams of a component S by $\mathbb{I}(S)$ (function $IN :: CSet \Rightarrow chanID\ set$ in Isabelle) and $\mathbb{O}(S)$ (function $OUT :: CSet \Rightarrow chanID\ set$ in Isabelle). The set of local variables of components is defined in Isabelle by VAR , and the function to map component identifiers to the corresponding variables is defined by $VAR :: CSet \Rightarrow varID\ set$.

Please note that concrete values for these functions cannot be specified in general, because they strongly depend on a concrete system. In this paper we present a small case study in the theories *DataDependenciesConcrete-Values.thy* (specification of the system architecture on several abstraction levels) and *DataDependenciesCaseStudy.thy* (proofs of system architectures' properties).

Function $subcomp :: CSet \Rightarrow CSet\ set$ maps components to a (possibly empty) set of its subcomponents.

We specify the components' dependencies by the function

$$Sources^L : CSet^L \rightarrow (CSet^L)^*$$

which returns for any component identifier A the corresponding (possibly empty) list of components (names) B_1, \dots, B_{AN} that are the sources for the input data streams of A (direct or indirect):

$$Sources^L(C) = DSources^L(C) \cup \bigcup_{S \in DSources^L(C)} \{S_1 \mid S_1 \in Sources^L(S)\}$$

Direct data dependencies are defined by the function

$$DSources^L : CSet^L \rightarrow (CSet^L)^*$$

$$DSources^L(C) = \{S \mid \exists x \in \mathbb{I}(C) \wedge x \in \mathbb{O}(S)\}$$

For example, $C_1 \in DSources^L(C_2)$ means that at least one of the output channels of C_1 is directly connected to some of input channels of C_2 .

$\mathbb{I}^D(C, y)$ denotes the subset of $\mathbb{I}(C)$ that output channel y depends upon, directly (specified in Isabelle by function $OUTfromCh:: chanID \Rightarrow chanID\ set$ or vial local variables (specified by function $OUTfromV:: chanID \Rightarrow varID\ set$). For example, let the values of the output channel y of component C depend only on the value of the local variable st that represents the current state of C and is updated depending to the input messages the component receives via the channel x , then $\mathbb{I}^D(C, y) = \{x\}$. In Isabelle, $\mathbb{I}^D(C, y)$ is specified by function $OUTfrom:: chanID \Rightarrow varID\ set$.

Based on the definition above, we can decompose system's components to have for each component's output channel the minimal subcomponent computing the corresponding results (we call them *elementary components*). An elementary component either

- should have a single output channel (in this case this component can have no local variables), or
- all its output channels are correlated, i.e. mutually depend on the same local variable(s).

If after this decomposition a single component is too complex, we can apply the decomposition strategy presented in [5].

For any component C , the dual function \mathbb{O}^D returns the corresponding set $\mathbb{O}^D(C, x)$ of output channels depending on input x . This is useful for tracing, e.g., if there are some changes in the specification, properties, constraints, etc. for x , we can trace which other channels can be affected by these changes.

If the input part of the component's interface is specified correctly in the sense that the component does not have any "unused" input channels, the following relation will hold: $\forall x \in \mathbb{I}(C). \mathbb{O}^D(C, x) \neq \emptyset$.

We illustrate the presented ideas by a small case study: we show how system's components can be decomposed to optimise the data dependencies within each single component, and after that we optimise architecture of the whole system. System S (cf. also Fig. 1) has 5 components, the set $CSet$ on the level L_0 is defined by $\{A_1, \dots, A_9\}$. The sets \mathbb{I}^D of data dependencies between the components are defined in the theory *DataDependenciesConcreteValues.thy*. We represent the dependencies graphically using dashed lines over the component box.

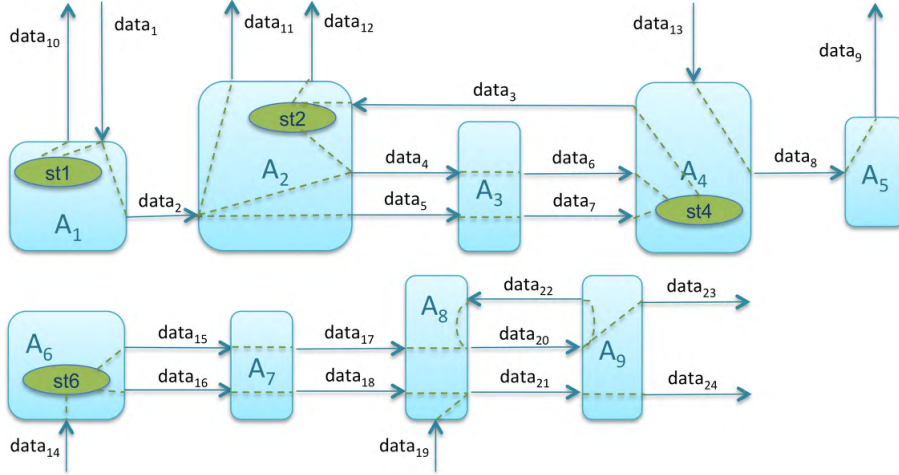


Figure 1: System S : Data dependencies and \mathbb{I}^D sets

Now we can decompose the system's components according to the given \mathbb{I}^D specification. This results into the next abstraction level L_1 of logical architecture (cf. Fig. 2), on which all components are elementary. Thus, we obtain a (flat) architecture of system. The main feature of this architecture is that each output channel (within the system) belongs the minimal sub-component of a system computing the corresponding results. We represent this (flat) architecture as a directed graph (components become vertices and channels become edges) and apply one of the existing distributed algorithms for the decomposition into its strongly connected components, e.g. FB [2], OBF [1], or the colouring algorithm [4]. Fig. 3 presents the result of the architecture optimisation.

After optimisation of system's architecture, we can find the minimal part of the system needed to check a specific property (cf. theory *DataDependencies*). A property can be represented by relations over data flows on the system's channels, and first of all we should check the property itself, whether it reflect a real relation within a system. Let for a relation r , I_r O_r be the sets of input and output channels of the system used in this relation. For each channel from O_r we recursively compute all the sets of the dependent components and corresponding input channels. Their union, restricted to the input channels of the system, should be equal to I_r , otherwise we should check whether the property was specified correctly.

Thus, from O_r we obtain the set *outSetOfComponents* of components having these channels as outputs, and compute the union of corresponding sources' sets. This union together with *outSetOfComponents* give us the minimal part of the system needed to check the property r : we formalise it in Isabelle by the predicate *minSetOfComponents*.

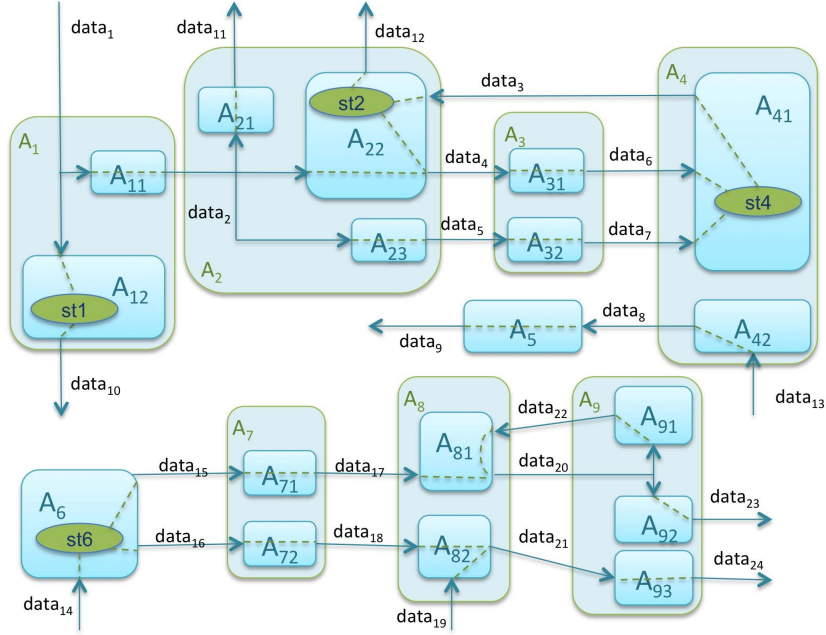


Figure 2: Components' decomposition (level L_1)

For each channel and elementary component (i.e. for any component on the abstraction level L_1) we specify the following measures:

- measure for costs of the data transfer/ upload to the cloud $UplSize(f)$: size of messages (data packages) within a data flow f and frequency they are produced. This measure can be defined on the level of logical modelling, where we already know the general type of the data and can also analyse the corresponding component (or environment) model to estimate the frequency the data are produced;
- measure for requirement of using high-performance computing and cloud virtual machines, $Perf(X)$: complexity of the computation within a component X , which can be estimated on the level of logical modelling as well.

On this basis, we build a system architecture, optimised for remote computation. The $UplSize$ measure should be analysed only for the channels that aren't local for the components on abstraction levels L_2 and L_3 .

Using graphical representation, we denote the channels with $UplSize$ measure higher than a predefined value by thick red arrows (cf. also set $UplSizeHighLoad$ in Isabelle theory $DataDependenciesConcreteValues.thy$), and the components with $Perf$ measure higher than a predefined value by light green colour (cf. also set $HighPerfSet$ in Isabelle theory $DataDependenciesConcreteValues.thy$), where all other channel and components are marked blue.

Fig. 4 represents a system architecture, optimised for remote computation: components from the abstraction level L_2 are composed together on the abstraction level L_3 , if they are connected by at least one channel with $UplSize$ measure higher than a predefined value. The components S'_4 and S'_7 have $Perf$ measure higher than a predefined value, i.e. using high-performance computing and cloud virtual machines is required.

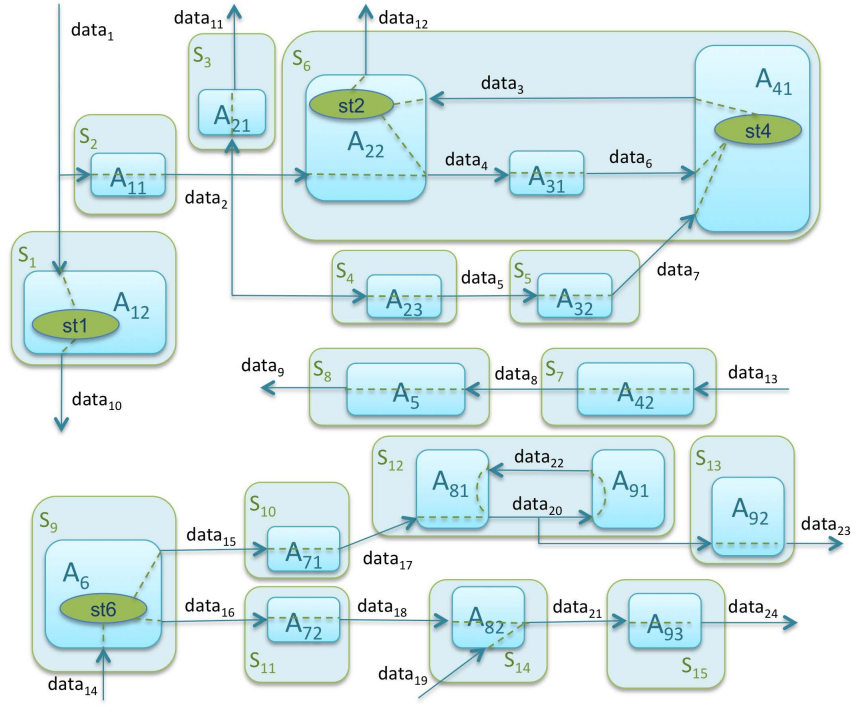


Figure 3: Architecture of S (level L_2)

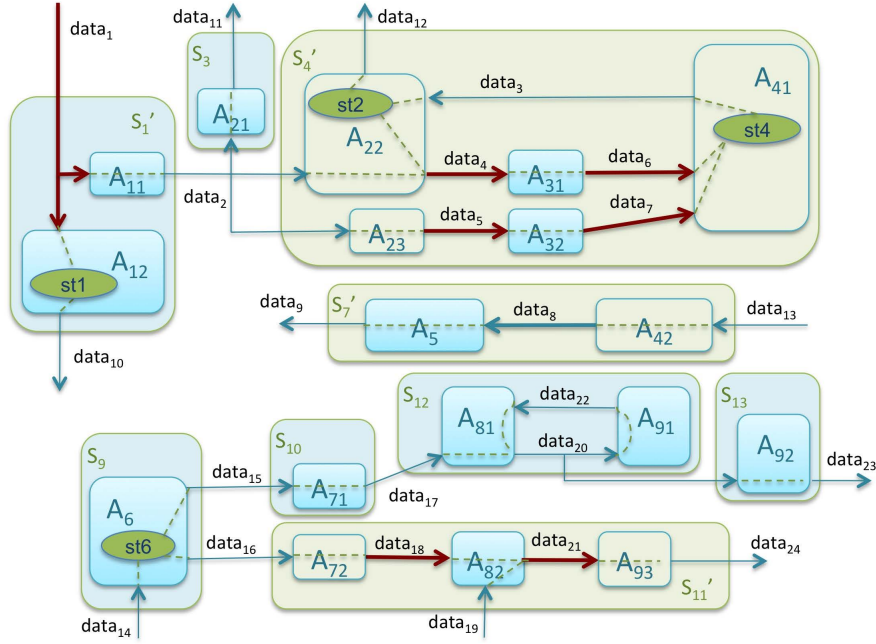


Figure 4: Optimised architecture of S (Level L_3)

2 Case Study: Definitions

```
theory DataDependenciesConcreteValues
```

```
  imports Main
```

```
begin
```

```
datatype CSet = sA1 | sA2 | sA3 | sA4 | sA5 | sA6 | sA7 | sA8 | sA9 |
              sA11 | sA12 | sA21 | sA22 | sA23 | sA31 | sA32 | sA41 | sA42 |
              sA71 | sA72 | sA81 | sA82 | sA91 | sA92 | sA93 |
              sS1 | sS2 | sS3 | sS4 | sS5 | sS6 | sS7 | sS8 | sS9 | sS10 | sS11 |
              sS12 | sS13 | sS14 | sS15 | sS1opt | sS4opt | sS7opt | sS11opt
```

```
datatype chanID = data1 | data2 | data3 | data4 | data5 | data6 | data7 |
                data8 | data9 | data10 | data11 | data12 | data13 | data14 | data15 |
                data16 | data17 | data18 | data19 | data20 | data21 | data22 | data23 | data24
```

```
datatype varID = stA1 | stA2 | stA4 | stA6
```

```
datatype AbstrLevelsID = level0 | level1 | level2 | level3
```

— function IN maps component ID to the set of its input channels

```
fun IN :: CSet  $\Rightarrow$  chanID set
```

```
where
```

```
  IN sA1 = { data1 }
```

```

| IN sA2 = { data2, data3 }
| IN sA3 = { data4, data5 }
| IN sA4 = { data6, data7, data13 }
| IN sA5 = { data8 }
| IN sA6 = { data14 }
| IN sA7 = { data15, data16 }
| IN sA8 = { data17, data18, data19, data22 }
| IN sA9 = { data20, data21 }
| IN sA11 = { data1 }
| IN sA12 = { data1 }
| IN sA21 = { data2 }
| IN sA22 = { data2, data3 }
| IN sA23 = { data2 }
| IN sA31 = { data4 }
| IN sA32 = { data5 }
| IN sA41 = { data6, data7 }
| IN sA42 = { data13 }
| IN sA71 = { data15 }
| IN sA72 = { data16 }
| IN sA81 = { data17, data22 }
| IN sA82 = { data18, data19 }
| IN sA91 = { data20 }
| IN sA92 = { data20 }
| IN sA93 = { data21 }
| IN sS1 = { data1 }
| IN sS2 = { data1 }
| IN sS3 = { data2 }
| IN sS4 = { data2 }
| IN sS5 = { data5 }
| IN sS6 = { data2, data7 }
| IN sS7 = { data13 }
| IN sS8 = { data8 }
| IN sS9 = { data14 }
| IN sS10 = { data15 }
| IN sS11 = { data16 }
| IN sS12 = { data17 }
| IN sS13 = { data20 }
| IN sS14 = { data18, data19 }
| IN sS15 = { data21 }
| IN sS1opt = { data1 }
| IN sS4opt = { data2 }
| IN sS7opt = { data13 }
| IN sS11opt = { data16, data19 }

```

— function `OUT` maps component ID to the set of its output channels

fun `OUT` :: `CSet` \Rightarrow `chanID set`

where

```

    OUT sA1 = { data2, data10 }
| OUT sA2 = { data4, data5, data11, data12 }

```



```

| OUT sA3 = { data6, data7 }
| OUT sA4 = { data3, data8 }
| OUT sA5 = { data9 }
| OUT sA6 = { data15, data16 }
| OUT sA7 = { data17, data18 }
| OUT sA8 = { data20, data21 }
| OUT sA9 = { data22, data23, data24 }
| OUT sA11 = { data2 }
| OUT sA12 = { data10 }
| OUT sA21 = { data11 }
| OUT sA22 = { data4, data12 }
| OUT sA23 = { data5 }
| OUT sA31 = { data6 }
| OUT sA32 = { data7 }
| OUT sA41 = { data3 }
| OUT sA42 = { data8 }
| OUT sA71 = { data17 }
| OUT sA72 = { data18 }
| OUT sA81 = { data20 }
| OUT sA82 = { data21 }
| OUT sA91 = { data22 }
| OUT sA92 = { data23 }
| OUT sA93 = { data24 }
| OUT sS1 = { data10 }
| OUT sS2 = { data2 }
| OUT sS3 = { data11 }
| OUT sS4 = { data5 }
| OUT sS5 = { data7 }
| OUT sS6 = { data12 }
| OUT sS7 = { data8 }
| OUT sS8 = { data9 }
| OUT sS9 = { data15, data16 }
| OUT sS10 = { data17 }
| OUT sS11 = { data18 }
| OUT sS12 = { data20 }
| OUT sS13 = { data23 }
| OUT sS14 = { data21 }
| OUT sS15 = { data24 }
| OUT sS1opt = { data2, data10 }
| OUT sS4opt = { data12 }
| OUT sS7opt = { data9 }
| OUT sS11opt = { data24 }

```

— function VAR maps component IDs to the set of its local variables

fun VAR :: CSet \Rightarrow varID set

where

```

    VAR sA1 = { stA1 }
| VAR sA2 = { stA2 }

```

```

| VAR sA3 = {}
| VAR sA4 = { stA4 }
| VAR sA5 = {}
| VAR sA6 = { stA6 }
| VAR sA7 = {}
| VAR sA8 = {}
| VAR sA9 = {}
| VAR sA11 = {}
| VAR sA12 = { stA1 }
| VAR sA21 = {}
| VAR sA22 = { stA2 }
| VAR sA23 = {}
| VAR sA31 = {}
| VAR sA32 = {}
| VAR sA41 = {stA4 }
| VAR sA42 = {}
| VAR sA71 = {}
| VAR sA72 = {}
| VAR sA81 = {}
| VAR sA82 = {}
| VAR sA91 = {}
| VAR sA92 = {}
| VAR sA93 = {}
| VAR sS1 = { stA1 }
| VAR sS2 = {}
| VAR sS3 = {}
| VAR sS4 = {}
| VAR sS5 = {}
| VAR sS6 = {stA2, stA4}
| VAR sS7 = {}
| VAR sS8 = {}
| VAR sS9 = {stA6}
| VAR sS10 = {}
| VAR sS11 = {}
| VAR sS12 = {}
| VAR sS13 = {}
| VAR sS14 = {}
| VAR sS15 = {}
| VAR sS1opt = { stA1 }
| VAR sS4opt = { stA2, stA4 }
| VAR sS7opt = {}
| VAR sS11opt = {}

```

— function `subcomp` maps component ID to the set of its subcomponents

fun `subcomp` :: *CSet* ⇒ *CSet set*

where

```

  subcomp sA1 = { sA11, sA12 }
| subcomp sA2 = { sA21, sA22, sA23 }

```

```

| subcomp sA3 = { sA31, sA32 }
| subcomp sA4 = { sA41, sA42 }
| subcomp sA5 = {}
| subcomp sA6 = {}
| subcomp sA7 = { sA71, sA72 }
| subcomp sA8 = { sA81, sA82 }
| subcomp sA9 = { sA91, sA92, sA93 }
| subcomp sA11 = {}
| subcomp sA12 = {}
| subcomp sA21 = {}
| subcomp sA22 = {}
| subcomp sA23 = {}
| subcomp sA31 = {}
| subcomp sA32 = {}
| subcomp sA41 = {}
| subcomp sA42 = {}
| subcomp sA71 = {}
| subcomp sA72 = {}
| subcomp sA81 = {}
| subcomp sA82 = {}
| subcomp sA91 = {}
| subcomp sA92 = {}
| subcomp sA93 = {}
| subcomp sS1 = { sA12 }
| subcomp sS2 = { sA11 }
| subcomp sS3 = { sA21 }
| subcomp sS4 = { sA23 }
| subcomp sS5 = { sA32 }
| subcomp sS6 = { sA22, sA31, sA41 }
| subcomp sS7 = { sA42 }
| subcomp sS8 = { sA5 }
| subcomp sS9 = { sA6 }
| subcomp sS10 = { sA71 }
| subcomp sS11 = { sA72 }
| subcomp sS12 = { sA81, sA91 }
| subcomp sS13 = { sA92 }
| subcomp sS14 = { sA82 }
| subcomp sS15 = { sA93 }
| subcomp sS1opt = { sA11, sA12 }
| subcomp sS4opt = { sA22, sA23, sA31, sA32, sA41 }
| subcomp sS7opt = { sA42, sA5 }
| subcomp sS11opt = { sA72, sA82, sA93 }

```

— function `AbstrLevel` maps abstraction level ID to the corresponding set of components

axiomatization

AbstrLevel :: *AbstrLevelsID* ⇒ *CSet set*

where

AbstrLevel0:

```

AbstrLevel level0 = {sA1, sA2, sA3, sA4, sA5, sA6, sA7, sA8, sA9}
and
AbstrLevel1:
AbstrLevel level1 = {sA11, sA12, sA21, sA22, sA23, sA31, sA32,
  sA41, sA42, sA5, sA6, sA71, sA72, sA81, sA82, sA91, sA92, sA93}
and
AbstrLevel2:
AbstrLevel level2 = {sS1, sS2, sS3, sS4, sS5, sS6, sS7, sS8,
  sS9, sS10, sS11, sS12, sS13, sS14, sS15}
and
AbstrLevel3:
AbstrLevel level3 = {sS1opt, sS3, sS4opt, sS7opt, sS9, sS10, sS11opt, sS12, sS13
}

```

— function VARfrom maps variable ID to the set of input channels it depends from

```

fun VARfrom :: varID ⇒ chanID set
where
  VARfrom stA1 = {data1}
| VARfrom stA2 = {data3}
| VARfrom stA4 = {data6, data7}
| VARfrom stA6 = {data14}

```

— function VARto maps variable ID to the set of output channels depending from this variable

```

fun VARto :: varID ⇒ chanID set
where
  VARto stA1 = {data10}
| VARto stA2 = {data4, data12}
| VARto stA4 = {data3}
| VARto stA6 = {data15, data16}

```

— function OUTfromCh maps channel ID to the set of input channels

— from which it depends directly;

— an empty set means that the channel is either input of the system or

— its values are computed from local variables or are generated

— within some component independently

```

fun OUTfromCh :: chanID ⇒ chanID set
where
  OUTfromCh data1 = {}
| OUTfromCh data2 = {data1}
| OUTfromCh data3 = {}
| OUTfromCh data4 = {data2}
| OUTfromCh data5 = {data2}
| OUTfromCh data6 = {data4}
| OUTfromCh data7 = {data5}
| OUTfromCh data8 = {data13}
| OUTfromCh data9 = {data8}
| OUTfromCh data10 = {}
| OUTfromCh data11 = {data2}

```

```

| OUTfromCh data12 = {}
| OUTfromCh data13 = {}
| OUTfromCh data14 = {}
| OUTfromCh data15 = {}
| OUTfromCh data16 = {}
| OUTfromCh data17 = {data15}
| OUTfromCh data18 = {data16}
| OUTfromCh data19 = {}
| OUTfromCh data20 = {data17, data22}
| OUTfromCh data21 = {data18, data19}
| OUTfromCh data22 = {data20}
| OUTfromCh data23 = {data21}
| OUTfromCh data24 = {data20}

```

— function *OUTfromV* maps channel ID to the set of local variables it depends from

```

fun OUTfromV :: chanID ⇒ varID set
where

```

```

  OUTfromV data1 = {}
| OUTfromV data2 = {}
| OUTfromV data3 = {stA4}
| OUTfromV data4 = {stA2}
| OUTfromV data5 = {}
| OUTfromV data6 = {}
| OUTfromV data7 = {}
| OUTfromV data8 = {}
| OUTfromV data9 = {}
| OUTfromV data10 = {stA1}
| OUTfromV data11 = {}
| OUTfromV data12 = {stA2}
| OUTfromV data13 = {}
| OUTfromV data14 = {}
| OUTfromV data15 = {stA6}
| OUTfromV data16 = {stA6}
| OUTfromV data17 = {}
| OUTfromV data18 = {}
| OUTfromV data19 = {}
| OUTfromV data20 = {}
| OUTfromV data21 = {}
| OUTfromV data22 = {}
| OUTfromV data23 = {}
| OUTfromV data24 = {}

```

— Set of channels which have *UplSize* measure greather that the predefined value *HighLoad*

definition

```

UplSizeHighLoad :: chanID set

```

where

```

UplSizeHighLoad ≡ {data1, data4, data5, data6, data7, data8, data18, data21}

```

— Set of components from the abstraction level 1 for which the Perf measure is greater than the predefined value *HighPerf*

definition

HighPerfSet :: CSet set

where

HighPerfSet ≡ {*sA22*, *sA23*, *sA41*, *sA42*, *sA72*, *sA93*}

end

3 Inter-/Intracomponent dependencies

theory *DataDependencies*

imports *DataDependenciesConcreteValues*

begin

— component and its subcomponents should be defined on different abstraction levels

definition

correctCompositionDiffLevels :: CSet ⇒ bool

where

correctCompositionDiffLevels *S* ≡

∀ *C* ∈ *subcomp* *S*. ∀ *i*. *S* ∈ *AbstrLevel* *i* → *C* ∉ *AbstrLevel* *i*

— General system's property: for all abstraction levels and all components should hold

— component and its subcomponents should be defined on different abstraction levels

definition

correctCompositionDiffLevelsSYSTEM :: bool

where

correctCompositionDiffLevelsSYSTEM ≡

(∀ *S*::CSet. (*correctCompositionDiffLevels* *S*))

— if a local variable belongs to one of the subcomponents, it also belongs to the composed component

definition

correctCompositionVAR :: CSet ⇒ bool

where

correctCompositionVAR *S* ≡

∀ *C* ∈ *subcomp* *S*. ∀ *v* ∈ *VAR* *C*. *v* ∈ *VAR* *S*

— General system's property: for all abstraction levels and all components should hold

— if a local variable belongs to one of the subcomponents, it also belongs to the composed component

definition

correctCompositionVARSYSTEM :: bool

where

$correctCompositionVARSYSTEM \equiv$
 $(\forall S::CSet. (correctCompositionVAR S))$

— after correct decomposition of a component each of its local variable can belong only to one of its subcomponents

definition

$correctDeCompositionVAR :: CSet \Rightarrow bool$

where

$correctDeCompositionVAR S \equiv$
 $\forall v \in VAR S. \forall C1 \in subcomp S. \forall C2 \in subcomp S. v \in VAR C1 \wedge v \in VAR C2 \longrightarrow C1 = C2$

— General system's property: for all abstraction levels and all components should hold

— after correct decomposition of a component each of its local variable can belong only to one of its subcomponents

definition

$correctDeCompositionVARSYSTEM :: bool$

where

$correctDeCompositionVARSYSTEM \equiv$
 $(\forall S::CSet. (correctDeCompositionVAR S))$

— if x is an output channel of a component C on some anstraction level, it cannot be an output of another component on the same level

definition

$correctCompositionOUT :: chanID \Rightarrow bool$

where

$correctCompositionOUT x \equiv$
 $\forall C i. x \in OUT C \wedge C \in AbstrLevel i \longrightarrow (\forall S \in AbstrLevel i. x \notin OUT S)$

— General system's property: for all abstraction levels and all channels should hold

definition

$correctCompositionOUTSYSTEM :: bool$

where

$correctCompositionOUTSYSTEM \equiv (\forall x. correctCompositionOUT x)$

— if X is a subcomponent of a component C on some anstraction level, it cannot be a subcomponent of another component on the same level

definition

$correctCompositionSubcomp :: CSet \Rightarrow bool$

where

$correctCompositionSubcomp X \equiv$
 $\forall C i. X \in subcomp C \wedge C \in AbstrLevel i \longrightarrow (\forall S \in AbstrLevel i. (S \neq C \longrightarrow X \notin subcomp S))$

— General system's property: for all abstraction levels and all components should hold

definition

$correctCompositionSubcompSYSTEM :: bool$

where

$correctCompositionSubcompSYSTEM \equiv (\forall X. correctCompositionSubcomp X)$

— If a component belongs is defined in the set CSet, it should belong to at least one abstraction level

definition

$allComponentsUsed :: bool$

where

$allComponentsUsed \equiv \forall C. \exists i. C \in AbstrLevel i$

— if a component does not have any local variables, none of its subcomponents has any local variables

lemma *correctDeCompositionVARempty:*

assumes $correctCompositionVAR S$

and $VAR S = \{\}$

shows $\forall C \in subcomp S. VAR C = \{\}$

using *assms* **by** (*metis all-not-in-conv correctCompositionVAR-def*)

— function OUTfrom maps channel ID to the set of input channels it depends from,

— directly (OUTfromCh) or via local variables (VARfrom)

— an empty set means that the channel is either input of the system or

— its values are generated within some component independently

definition $OUTfrom :: chanID \Rightarrow chanID\ set$

where

$OUTfrom x \equiv (OUTfromCh x) \cup \{y. \exists v. v \in (OUTfromV x) \wedge y \in (VARfrom v)\}$

— if x depends from some input channel(s) directly, then exists

— a component which has them as input channels and x as an output channel

definition

$OUTfromChCorrect :: chanID \Rightarrow bool$

where

$OUTfromChCorrect x \equiv$

$(OUTfromCh x \neq \{\}) \longrightarrow$

$(\exists Z. (x \in (OUT Z) \wedge (\forall y \in (OUTfromCh x). y \in IN Z)))$

— General system's property: for channels in the system should hold:

— if x depends from some input channel(s) directly, then exists

— a component which has them as input channels and x as an output channel

definition

$OUTfromChCorrectSYSTEM :: bool$

where

$OUTfromChCorrectSYSTEM \equiv (\forall x::chanID. (OUTfromChCorrect x))$

— if x depends from some local variables, then exists a component

— to which these variables belong and which has x as an output channel

definition

$OUTfromVCorrect1 :: chanID \Rightarrow bool$

where

$OUTfromVCorrect1 x \equiv$

$(OUTfromV x \neq \{\}) \longrightarrow$

$(\exists Z . (x \in (OUT Z) \wedge (\forall v \in (OUTfromV x). v \in VAR Z)))$

— General system's property: for channels in the system should hold the above property:

definition

$OUTfromVCorrect1SYSTEM :: bool$

where

$OUTfromVCorrect1SYSTEM \equiv (\forall x::chanID. (OUTfromVCorrect1 x))$

— if x does not depend from any local variables, then it does not belong to any set VARfrom

definition

$OUTfromVCorrect2 :: chanID \Rightarrow bool$

where

$OUTfromVCorrect2 x \equiv$

$(OUTfromV x = \{\}) \longrightarrow (\forall v::varID. x \notin (VARto v))$

— General system's property: for channels in the system should hold the above property:

definition

$OUTfromVCorrect2SYSTEM :: bool$

where

$OUTfromVCorrect2SYSTEM \equiv (\forall x::chanID. (OUTfromVCorrect2 x))$

— General system's property:

— definitions OUTfromV and VARto should give equivalent mappings

definition

$OUTfromV-VARto :: bool$

where

$OUTfromV-VARto \equiv$

$(\forall x::chanID. \forall v::varID. (v \in OUTfromV x \longleftrightarrow x \in (VARto v)))$

— General system's property for abstraction levels 0 and 1

— if a variable v belongs to a component, then all the channels v

— depends from should be input channels of this component

definition

$VARfromCorrectSYSTEM :: bool$

where

$VARfromCorrectSYSTEM \equiv$

$(\forall v::varID. \forall Z \in ((AbstrLevel level0) \cup (AbstrLevel level1)).$

$(v \in VAR Z) \longrightarrow (\forall x \in VARfrom v. x \in IN Z))$

— General system's property for abstraction levels 0 and 1

— if a variable v belongs to a component, then all the channels v

— provides value to should be input channels of this component

definition

$$VARtoCorrectSYSTEM :: bool$$
where

$$\begin{aligned} VARtoCorrectSYSTEM &\equiv \\ &(\forall v::varID. \forall Z \in ((AbstrLevel level0) \cup (AbstrLevel level1)). \\ &((v \in VAR Z) \longrightarrow (\forall x \in VARto v. x \in OUT Z))) \end{aligned}$$

— to detect local variables, unused for computation of any output

definition

$$VARusefulSYSTEM :: bool$$
where

$$VARusefulSYSTEM \equiv (\forall v::varID. (VARto v \neq \{\}))$$
lemma

$$OUTfromV-VARto-lemma:$$

assumes $OUTfromV x \neq \{\}$ **and** $OUTfromV-VARto$

shows $\exists v::varID. x \in (VARto v)$

using *assms* **by** (*simp add: OUTfromV-VARto-def, auto*)

3.1 Direct and indirect data dependencies between components

— The component C should be defined on the same abstraction

— level we are searching for its direct or indirect sources,

— otherwise we get an empty set as result

definition

$$DSources :: AbstrLevelsID \Rightarrow CSet \Rightarrow CSet set$$
where

$$DSources i C \equiv \{Z. \exists x. x \in (IN C) \wedge x \in (OUT Z) \wedge Z \in (AbstrLevel i) \wedge C \in (AbstrLevel i)\}$$
lemma $DSourcesLevelX:$

$$(DSources i X) \subseteq (AbstrLevel i)$$

by (*simp add: DSources-def, auto*)

— The component C should be defined on the same abstraction level we are

— searching for its direct or indirect acceptors (components, for which C is a source),

— otherwise we get an empty set as result

definition

$$DAcc :: AbstrLevelsID \Rightarrow CSet \Rightarrow CSet set$$
where

$$DAcc i C \equiv \{Z. \exists x. x \in (OUT C) \wedge x \in (IN Z) \wedge Z \in (AbstrLevel i) \wedge C \in (AbstrLevel i)\}$$
axiomatization

$$Sources :: AbstrLevelsID \Rightarrow CSet \Rightarrow CSet set$$
where

$$SourcesDef:$$

$(Sources\ i\ C) = (DSources\ i\ C) \cup (\bigcup S \in (DSources\ i\ C). (Sources\ i\ S))$
and
SourceExistsDSource:
 $S \in (Sources\ i\ C) \longrightarrow (\exists Z. S \in (DSources\ i\ Z))$
and
NDSourceExistsDSource:
 $S \in (Sources\ i\ C) \wedge S \notin (DSources\ i\ C) \longrightarrow$
 $(\exists Z. S \in (DSources\ i\ Z) \wedge Z \in (Sources\ i\ C))$
and
SourcesTrans:
 $(C \in Sources\ i\ S \wedge S \in Sources\ i\ Z) \longrightarrow C \in Sources\ i\ Z$
and
SourcesLevelX:
 $(Sources\ i\ X) \subseteq (AbstrLevel\ i)$
and
SourcesLoop:
 $(Sources\ i\ C) = (XS \cup (Sources\ i\ S)) \wedge (Sources\ i\ S) = (ZS \cup (Sources\ i\ C))$
 $\longrightarrow (Sources\ i\ C) = XS \cup ZS \cup \{C, S\}$
— if we have a loop in the dependencies we need to cut it for counting the sources

axiomatization

$Acc :: AbstrLevelsID \Rightarrow CSet \Rightarrow CSet\ set$
where
AccDef:
 $(Acc\ i\ C) = (DAcc\ i\ C) \cup (\bigcup S \in (DAcc\ i\ C). (Acc\ i\ S))$
and
Acc-Sources:
 $(X \in Acc\ i\ C) = (C \in Sources\ i\ X)$
and
AccSngleLoop:
 $DAcc\ i\ C = \{S\} \wedge DAcc\ i\ S = \{C\} \longrightarrow Acc\ i\ C = \{C, S\}$
and
AccLoop:
 $(Acc\ i\ C) = (XS \cup (Acc\ i\ S)) \wedge (Acc\ i\ S) = (ZS \cup (Acc\ i\ C))$
 $\longrightarrow (Acc\ i\ C) = XS \cup ZS \cup \{C, S\}$
— if we have a loop in the dependencies we need to cut it for counting the accessors

lemma *Acc-SourcesNOT:* $(X \notin Acc\ i\ C) = (C \notin Sources\ i\ X)$
by (*metis Acc-Sources*)

— component S is not a source for any component on the abstraction level i

definition

$isNotDSource :: AbstrLevelsID \Rightarrow CSet \Rightarrow bool$

where

$isNotDSource\ i\ S \equiv (\forall x \in (OUT\ S). (\forall Z \in (AbstrLevel\ i). (x \notin (IN\ Z))))$

— component S is not a source for a component Z on the abstraction level i

definition

$isNotDSourceX :: AbstrLevelsID \Rightarrow CSet \Rightarrow CSet \Rightarrow bool$

where

$isNotDSourceX\ i\ S\ C \equiv (\forall x \in (OUT\ S). (C \notin (AbstrLevel\ i) \vee (x \notin (IN\ C))))$

lemma *isNotSource-isNotSourceX*:

$isNotDSource\ i\ S = (\forall C. isNotDSourceX\ i\ S\ C)$

by (*auto*, (*simp add: isNotDSource-def isNotDSourceX-def*))+)

lemma *DAcc-DSources*:

$(X \in DAcc\ i\ C) = (C \in DSources\ i\ X)$

by (*auto*, (*simp add: DAcc-def DSources-def, auto*))+)

lemma *DAcc-DSourcesNOT*:

$(X \notin DAcc\ i\ C) = (C \notin DSources\ i\ X)$

by (*auto*, (*simp add: DAcc-def DSources-def, auto*))+)

lemma *DSource-level*:

assumes $S \in (DSources\ i\ C)$

shows $C \in (AbstrLevel\ i)$

using *assms* **by** (*simp add: DSources-def, auto*)

lemma *SourceExistsDSource-level*:

assumes $S \in (Sources\ i\ C)$

shows $\exists Z \in (AbstrLevel\ i). (S \in (DSources\ i\ Z))$

using *assms* **by** (*metis DSource-level SourceExistsDSource*)

lemma *Sources-DSources*:

$(DSources\ i\ C) \subseteq (Sources\ i\ C)$

proof –

have $(Sources\ i\ C) = (DSources\ i\ C) \cup (\bigcup S \in (DSources\ i\ C). (Sources\ i\ S))$

by (*rule SourcesDef*)

thus *?thesis* **by** *auto*

qed

lemma *NoDSourceNoSource*:

assumes $S \notin (Sources\ i\ C)$

shows $S \notin (DSources\ i\ C)$

using *assms* **by** (*metis (full-types) Sources-DSources rev-subsetD*)

lemma *DSourcesEmptySources*:

assumes $DSources\ i\ C = \{\}$

shows $Sources\ i\ C = \{\}$

proof –

have $(Sources\ i\ C) = (DSources\ i\ C) \cup (\bigcup S \in (DSources\ i\ C). (Sources\ i\ S))$

by (*rule SourcesDef*)

with *assms* **show** *?thesis* **by** *auto*

qed

lemma *DSource-Sources*:

assumes $S \in (DSources\ i\ C)$

shows $(Sources\ i\ S) \subseteq (Sources\ i\ C)$
proof –
have $(Sources\ i\ C) = (DSources\ i\ C) \cup (\bigcup S \in (DSources\ i\ C). (Sources\ i\ S))$
by *(rule SourcesDef)*
with *assms* **show** *?thesis* **by** *auto*
qed

lemma *SourcesOnlyDSources*:
assumes $\forall X. (X \in (DSources\ i\ C) \longrightarrow (DSources\ i\ X) = \{\})$
shows $Sources\ i\ C = DSources\ i\ C$
proof –
have *sDef*: $(Sources\ i\ C) = (DSources\ i\ C) \cup (\bigcup S \in (DSources\ i\ C). (Sources\ i\ S))$
by *(rule SourcesDef)*
from *assms* **have** $\forall X. (X \in (DSources\ i\ C) \longrightarrow (Sources\ i\ X) = \{\})$
by *(simp add: DSourcesEmptySources)*
hence $(\bigcup S \in (DSources\ i\ C). (Sources\ i\ S)) = \{\}$ **by** *auto*
with *sDef* **show** *?thesis* **by** *simp*
qed

lemma *SourcesEmptyDSources*:
assumes $Sources\ i\ C = \{\}$
shows $DSources\ i\ C = \{\}$
using *assms* **by** *(metis Sources-DSources bot.extremum-uniqueI)*

lemma *NotDSource*:
assumes $\forall x \in (OUT\ S). (\forall Z \in (AbstrLevel\ i). (x \notin (IN\ Z)))$
shows $\forall C \in (AbstrLevel\ i). S \notin (DSources\ i\ C)$
using *assms* **by** *(simp add: AbstrLevel0 DSources-def)*

lemma *allNotDSource-NotSource*:
assumes $\forall C. S \notin (DSources\ i\ C)$
shows $\forall Z. S \notin (Sources\ i\ Z)$
using *assms* **by** *(metis SourceExistsDSource)*

lemma *NotDSource-NotSource*:
assumes $\forall C \in (AbstrLevel\ i). S \notin (DSources\ i\ C)$
shows $\forall Z \in (AbstrLevel\ i). S \notin (Sources\ i\ Z)$
using *assms* **by** *(metis SourceExistsDSource-level)*

lemma *isNotSource-Sources*:
assumes *isNotDSource i S*
shows $\forall C \in (AbstrLevel\ i). S \notin (Sources\ i\ C)$
using *assms*
by *(simp add: isNotDSource-def, metis (full-types) NotDSource NotDSource-NotSource)*

lemma *SourcesAbstrLevel*:
assumes $x \in Sources\ i\ S$
shows $x \in AbstrLevel\ i$

using *assms*
by (*metis SourcesLevelX in-mono*)

lemma *DSourceIsSource*:

assumes $C \in DSources\ i\ S$
shows $C \in Sources\ i\ S$

proof –

have $(Sources\ i\ S) = (DSources\ i\ S) \cup (\bigcup Z \in (DSources\ i\ S). (Sources\ i\ Z))$
by (*rule SourcesDef*)

with *assms* **show** *?thesis* **by** *simp*

qed

lemma *DSourceOfDSource*:

assumes $Z \in DSources\ i\ S$
and $S \in DSources\ i\ C$
shows $Z \in Sources\ i\ C$

using *assms*

proof –

from *assms* **have** $src:Sources\ i\ S \subseteq Sources\ i\ C$ **by** (*simp add: DSource-Sources*)

from *assms* **have** $Z \in Sources\ i\ S$ **by** (*simp add: DSourceIsSource*)

with *src* **show** *?thesis* **by** *auto*

qed

lemma *SourceOfDSource*:

assumes $Z \in Sources\ i\ S$
and $S \in DSources\ i\ C$
shows $Z \in Sources\ i\ C$

using *assms*

proof –

from *assms* **have** $Sources\ i\ S \subseteq Sources\ i\ C$ **by** (*simp add: DSource-Sources*)

thus *?thesis* **by** (*metis (full-types) assms(1) rev-subsetD*)

qed

lemma *DSourceOfSource*:

assumes $cDS:C \in DSources\ i\ S$
and $sS:S \in Sources\ i\ Z$
shows $C \in Sources\ i\ Z$

proof –

from *cDS* **have** $C \in Sources\ i\ S$ **by** (*simp add: DSourceIsSource*)

from *this* **and** *sS* **show** *?thesis* **by** (*metis (full-types) SourcesTrans*)

qed

lemma *Sources-singleDSource*:

assumes $DSources\ i\ S = \{C\}$
shows $Sources\ i\ S = \{C\} \cup Sources\ i\ C$

proof –

have *sDef*: $(Sources\ i\ S) = (DSources\ i\ S) \cup (\bigcup Z \in (DSources\ i\ S). (Sources\ i\ Z))$

by (*rule SourcesDef*)

from *assms* **have** $(\bigcup Z \in (DSources\ i\ S)). (Sources\ i\ Z) = Sources\ i\ C$
by *auto*
with *sDef* *assms* **show** *?thesis* **by** *simp*
qed

lemma *Sources-2DSources*:

assumes $DSources\ i\ S = \{C1, C2\}$
shows $Sources\ i\ S = \{C1, C2\} \cup Sources\ i\ C1 \cup Sources\ i\ C2$
proof –
have *sDef*: $(Sources\ i\ S) = (DSources\ i\ S) \cup (\bigcup Z \in (DSources\ i\ S)). (Sources\ i\ Z)$
by (*rule SourcesDef*)
from *assms* **have** $(\bigcup Z \in (DSources\ i\ S)). (Sources\ i\ Z) = Sources\ i\ C1 \cup Sources\ i\ C2$
by *auto*
with *sDef* **and** *assms* **show** *?thesis* **by** *simp*
qed

lemma *Sources-3DSources*:

assumes $DSources\ i\ S = \{C1, C2, C3\}$
shows $Sources\ i\ S = \{C1, C2, C3\} \cup Sources\ i\ C1 \cup Sources\ i\ C2 \cup Sources\ i\ C3$
proof –
have *sDef*: $(Sources\ i\ S) = (DSources\ i\ S) \cup (\bigcup Z \in (DSources\ i\ S)). (Sources\ i\ Z)$
by (*rule SourcesDef*)
from *assms* **have** $(\bigcup Z \in (DSources\ i\ S)). (Sources\ i\ Z) = Sources\ i\ C1 \cup Sources\ i\ C2 \cup Sources\ i\ C3$
by *auto*
with *sDef* **and** *assms* **show** *?thesis* **by** *simp*
qed

lemma *singleDSourceEmpty4isNotDSource*:

assumes $DAcc\ i\ C = \{S\}$
and $Z \neq S$
shows $C \notin (DSources\ i\ Z)$
proof –
from *assms* **have** $(Z \notin DAcc\ i\ C)$ **by** *simp*
thus *?thesis* **by** (*simp add: DAcc-DSourcesNOT*)
qed

lemma *singleDSourceEmpty4isNotDSourceLevel*:

assumes $DAcc\ i\ C = \{S\}$
shows $\forall Z \in (AbstrLevel\ i). Z \neq S \longrightarrow C \notin (DSources\ i\ Z)$
using *assms* **by** (*metis singleDSourceEmpty4isNotDSource*)

lemma *isNotDSource-EmptyDAcc*:

assumes *isNotDSource* *i* *S*

shows $DAcc\ i\ S = \{\}$
using *assms* **by** (*simp add: DAcc-def isNotDSource-def, auto*)

lemma *isNotDSource-EmptyAcc:*

assumes *isNotDSource i S*

shows $Acc\ i\ S = \{\}$

proof –

have $(Acc\ i\ S) = (DAcc\ i\ S) \cup (\bigcup X \in (DAcc\ i\ S). (Acc\ i\ X))$

by (*rule AccDef*)

thus *?thesis* **by** (*metis SUP-empty Un-absorb assms isNotDSource-EmptyDAcc*)

qed

lemma *singleDSourceEmpty-Acc:*

assumes $DAcc\ i\ C = \{S\}$

and *isNotDSource i S*

shows $Acc\ i\ C = \{S\}$

proof –

have $AccC:(Acc\ i\ C) = (DAcc\ i\ C) \cup (\bigcup S \in (DAcc\ i\ C). (Acc\ i\ S))$

by (*rule AccDef*)

from *assms* **have** $Acc\ i\ S = \{\}$ **by** (*simp add: isNotDSource-EmptyAcc*)

with *AccC* **show** *?thesis*

by (*metis SUP-empty UN-insert Un-commute Un-empty-left assms(1)*)

qed

lemma *singleDSourceEmpty4isNotSource:*

assumes $DAcc\ i\ C = \{S\}$

and *nSourceS:isNotDSource i S*

and $Z \neq S$

shows $C \notin (Sources\ i\ Z)$

proof –

from *assms* **have** $Acc\ i\ C = \{S\}$ **by** (*simp add: singleDSourceEmpty-Acc*)

with *assms* **have** $Z \notin Acc\ i\ C$ **by** *simp*

thus *?thesis* **by** (*simp add: Acc-SourcesNOT*)

qed

lemma *singleDSourceEmpty4isNotSourceLevel:*

assumes $DAcc\ i\ C = \{S\}$

and *nSourceS:isNotDSource i S*

shows $\forall Z \in (AbstrLevel\ i). Z \neq S \longrightarrow C \notin (Sources\ i\ Z)$

using *assms*

by (*metis singleDSourceEmpty4isNotSource*)

lemma *singleDSourceLoop:*

assumes $DAcc\ i\ C = \{S\}$

and $DAcc\ i\ S = \{C\}$

shows $\forall Z \in (AbstrLevel\ i). (Z \neq S \wedge Z \neq C \longrightarrow C \notin (Sources\ i\ Z))$

using *assms*

by (*metis AccSigleLoop Acc-SourcesNOT empty-iff insert-iff*)

3.2 Components that are elementary wrt. data dependencies

— two output channels of a component C are corelated, if they mutually depend on the same local variable(s)

definition

$outPairCorelated :: CSet \Rightarrow chanID \Rightarrow chanID \Rightarrow bool$

where

$outPairCorelated C x y \equiv$
 $(x \in OUT C) \wedge (y \in OUT C) \wedge$
 $(OUTfromV x) \cap (OUTfromV y) \neq \{\}$

— We call a set of output channels of a component correlated to its output channel x ,

— if they mutually depend on the same local variable(s)

definition

$outSetCorelated :: chanID \Rightarrow chanID set$

where

$outSetCorelated x \equiv$
 $\{ y :: chanID . \exists v :: varID. (v \in (OUTfromV x) \wedge (y \in VARto v)) \}$

— Elementary component according to the data dependencies.

— This constraint should hold for all components on the abstraction level 1

definition

$elementaryCompDD :: CSet \Rightarrow bool$

where

$elementaryCompDD C \equiv$
 $((\exists x. (OUT C) = \{x\}) \vee$
 $(\forall x \in (OUT C). \forall y \in (OUT C). ((outSetCorelated x) \cap (outSetCorelated y)$
 $\neq \{\})))$

— the set $(outSetCorelated x)$ is empty if x does not depend from any variable

lemma $outSetCorelatedEmpty1$:

assumes $OUTfromV x = \{\}$

shows $outSetCorelated x = \{\}$

using $assms$ **by** ($simp$ $add: outSetCorelated-def$)

— if x depends from at least one variable and the predicates $OUTfromV$ and $VARto$ are defined correctly,

— the set $(outSetCorelated x)$ contains x itself

lemma $outSetCorelatedNonemptyX$:

assumes $OUTfromV x \neq \{\}$ **and** $correct3: OUTfromV-VARto$

shows $x \in outSetCorelated x$

proof —

from $assms$ **have** $\exists v :: varID. x \in (VARto v)$

by ($rule$ $OUTfromV-VARto-lemma$)

from $this$ **and** $assms$ **show** $?thesis$

by (*simp add: outSetCorelated-def OUTfromV-VARto-def*)
qed

— if the set (`outSetCorelated x`) is empty, this means that `x` does not depend from any variable

lemma *outSetCorelatedEmpty2*:
assumes *outSetCorelated x = {}* **and** *correct3:OUTfromV-VARto*
shows *OUTfromV x = {}*
proof (*rule ccontr*)
assume *OUTfromVNonempty:OUTfromV x ≠ {}*
from *this* **and** *correct3* **have** *x ∈ outSetCorelated x*
by (*rule outSetCorelatedNonemptyX*)
from *this* **and** *assms* **show** *False* **by** *simp*
qed

3.3 Set of components needed to check a specific property

— set of components specified on abstraction level `i`, which input channels belong to the set `chSet`

definition
inSetOfComponents :: *AbstrLevelsID* ⇒ *chanID set* ⇒ *CSet set*
where
inSetOfComponents i chSet ≡
 $\{X. (((IN X) \cap chSet \neq \{\}) \wedge X \in (AbstrLevel i))\}$

— Set of components from the abstraction level `i`, which output channels belong to the set `chSet`

definition
outSetOfComponents :: *AbstrLevelsID* ⇒ *chanID set* ⇒ *CSet set*
where
outSetOfComponents i chSet ≡
 $\{Y. (((OUT Y) \cap chSet \neq \{\}) \wedge Y \in (AbstrLevel i))\}$

— Set of components from the abstraction level `i`,
— which have output channels from the set `chSet` or are sources for such components

definition
minSetOfComponents :: *AbstrLevelsID* ⇒ *chanID set* ⇒ *CSet set*
where
minSetOfComponents i chSet ≡
 $(outSetOfComponents i chSet) \cup$
 $(\bigcup S \in (outSetOfComponents i chSet). (Sources i S))$

— Please note that a system output cannot beat the same time a local channel.

— channel `x` is a system input on an abstraction level `i`
definition *systemIN* :: *chanID* ⇒ *AbstrLevelsID* ⇒ *bool*
where

systemIN x i ≡ $(\exists C1 \in (AbstrLevel i). x \in (IN C1)) \wedge (\forall C2 \in (AbstrLevel i).$

$x \notin (OUT\ C2)$

— channel x is a system input on an abstraction level i

definition $systemOUT :: chanID \Rightarrow AbstrLevelsID \Rightarrow bool$

where

$systemOUT\ x\ i \equiv (\forall\ C1 \in (AbstrLevel\ i). x \notin (IN\ C1)) \wedge (\exists\ C2 \in (AbstrLevel\ i). x \in (OUT\ C2))$

— channel x is a system local channel on an abstraction level i

definition $systemLOC :: chanID \Rightarrow AbstrLevelsID \Rightarrow bool$

where

$systemLOC\ x\ i \equiv (\exists\ C1 \in (AbstrLevel\ i). x \in (IN\ C1)) \wedge (\exists\ C2 \in (AbstrLevel\ i). x \in (OUT\ C2))$

lemma $systemIN-noOUT$:

assumes $systemIN\ x\ i$

shows $\neg systemOUT\ x\ i$

using $assms$ **by** ($simp$ add : $systemIN-def\ systemOUT-def$)

lemma $systemOUT-noIN$:

assumes $systemOUT\ x\ i$

shows $\neg systemIN\ x\ i$

using $assms$ **by** ($simp$ add : $systemIN-def\ systemOUT-def$)

lemma $systemIN-noLOC$:

assumes $systemIN\ x\ i$

shows $\neg systemLOC\ x\ i$

using $assms$ **by** ($simp$ add : $systemIN-def\ systemLOC-def$)

lemma $systemLOC-noIN$:

assumes $systemLOC\ x\ i$

shows $\neg systemIN\ x\ i$

using $assms$ **by** ($simp$ add : $systemIN-def\ systemLOC-def$)

lemma $systemOUT-noLOC$:

assumes $systemOUT\ x\ i$

shows $\neg systemLOC\ x\ i$

using $assms$ **by** ($simp$ add : $systemOUT-def\ systemLOC-def$)

lemma $systemLOC-noOUT$:

assumes $systemLOC\ x\ i$

shows $\neg systemOUT\ x\ i$

using $assms$ **by** ($simp$ add : $systemLOC-def\ systemOUT-def$)

definition

$noIrrelevantChannels :: AbstrLevelsID \Rightarrow chanID\ set \Rightarrow bool$

where

$noIrrelevantChannels\ i\ chSet \equiv$

$\forall\ x \in chSet. ((systemIN\ x\ i) \longrightarrow$

$(\exists Z \in (\text{minSetOfComponents } i \text{ chSet}). x \in (\text{IN } Z))$

definition

$\text{allNeededINChannels} :: \text{AbstrLevelsID} \Rightarrow \text{chanID set} \Rightarrow \text{bool}$

where

$\text{allNeededINChannels } i \text{ chSet} \equiv$

$(\forall Z \in (\text{minSetOfComponents } i \text{ chSet}). \exists x \in (\text{IN } Z). ((\text{systemIN } x \ i) \longrightarrow (x \in \text{chSet})))$

— the set $(\text{outSetOfComponents } i \text{ chSet})$ should be a subset of all components specified on the abstraction level i

lemma $\text{outSetOfComponentsLimit}$:

$\text{outSetOfComponents } i \text{ chSet} \subseteq \text{AbstrLevel } i$

by $(\text{metis (lifting) mem-Collect-eq outSetOfComponents-def subsetI})$

— the set $(\text{inSetOfComponents } i \text{ chSet})$ should be a subset of all components specified on the abstraction level i

lemma $\text{inSetOfComponentsLimit}$:

$\text{inSetOfComponents } i \text{ chSet} \subseteq \text{AbstrLevel } i$

by $(\text{metis (lifting) inSetOfComponents-def mem-Collect-eq subsetI})$

— the set of components, which are sources for the components

— out of $(\text{inSetOfComponents } i \text{ chSet})$, should be a subset of

— all components specified on the abstraction level i

lemma SourcesLevelLimit :

$(\bigcup S \in (\text{outSetOfComponents } i \text{ chSet}). (\text{Sources } i \ S)) \subseteq \text{AbstrLevel } i$

proof –

have $\text{sg1: outSetOfComponents } i \text{ chSet} \subseteq \text{AbstrLevel } i$

by $(\text{simp add: outSetOfComponentsLimit})$

have $\forall S. S \in (\text{outSetOfComponents } i \text{ chSet}) \longrightarrow \text{Sources } i \ S \subseteq \text{AbstrLevel } i$

by $(\text{metis SourcesLevelX})$

from this and sg1 show ?thesis by auto

qed

lemma $\text{minSetOfComponentsLimit}$:

$\text{minSetOfComponents } i \text{ chSet} \subseteq \text{AbstrLevel } i$

proof –

have $\text{sg1: outSetOfComponents } i \text{ chSet} \subseteq \text{AbstrLevel } i$

by $(\text{simp add: outSetOfComponentsLimit})$

have $(\bigcup S \in (\text{outSetOfComponents } i \text{ chSet}). (\text{Sources } i \ S)) \subseteq \text{AbstrLevel } i$

by $(\text{simp add: SourcesLevelLimit})$

with sg1 show ?thesis by (simp add: minSetOfComponents-def)

qed

3.4 Additional properties: Remote Computation

— The value of $\text{UplSizeHighLoad } x$ is True if its UplSize measure is greater than

a predefined value

definition *UplSizeHighLoadCh* :: *chanID* \Rightarrow *bool*

where

UplSizeHighLoadCh *x* \equiv (*x* \in *UplSizeHighLoad*)

— if the *Perf* measure of at least one subcomponent is greater than a predefined value,

— the *Perf* measure of this component is greater than *HighPerf* too

axiomatization *HighPerfComp* :: *CSet* \Rightarrow *bool*

where

HighPerfComDef:

HighPerfComp *C* =

$((C \in \text{HighPerfSet}) \vee (\exists Z \in \text{subcomp } C. (\text{HighPerfComp } Z)))$

end

4 Case Study: Verification of Properties

theory *DataDependenciesCaseStudy*

imports *DataDependencies*

begin

4.1 Correct composition of components

— the lemmas *AbstrLevels* X Y with corresponding proofs can be composed

— and proven automatically, their proofs are identical

lemma *AbstrLevels-A1-A11*:

assumes *sA1* \in *AbstrLevel* *i*

shows *sA11* \notin *AbstrLevel* *i*

using *assms*

by (*induct* *i*, *simp* *add*: *AbstrLevel0*, *simp* *add*: *AbstrLevel1*, *simp* *add*: *AbstrLevel2*, *simp* *add*: *AbstrLevel3*)

lemma *AbstrLevels-A1-A12*:

assumes *sA1* \in *AbstrLevel* *i*

shows *sA12* \notin *AbstrLevel* *i*

lemma *AbstrLevels-A2-A21*:

assumes *sA2* \in *AbstrLevel* *i*

shows *sA21* \notin *AbstrLevel* *i*

lemma *AbstrLevels-A2-A22*:

assumes *sA2* \in *AbstrLevel* *i*

shows *sA22* \notin *AbstrLevel* *i*

lemma *AbstrLevels-A2-A23*:

assumes *sA2* \in *AbstrLevel* *i*

shows *sA23* \notin *AbstrLevel* *i*

lemma *AbstrLevels-A3-A31*:
assumes $sA3 \in \text{AbstrLevel } i$
shows $sA31 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A3-A32*:
assumes $sA3 \in \text{AbstrLevel } i$
shows $sA32 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A4-A41*:
assumes $sA4 \in \text{AbstrLevel } i$
shows $sA41 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A4-A42*:
assumes $sA4 \in \text{AbstrLevel } i$
shows $sA42 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A7-A71*:
assumes $sA7 \in \text{AbstrLevel } i$
shows $sA71 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A7-A72*:
assumes $sA7 \in \text{AbstrLevel } i$
shows $sA72 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A8-A81*:
assumes $sA8 \in \text{AbstrLevel } i$
shows $sA81 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A8-A82*:
assumes $sA8 \in \text{AbstrLevel } i$
shows $sA82 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A9-A91*:
assumes $sA9 \in \text{AbstrLevel } i$
shows $sA91 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A9-A92*:
assumes $sA9 \in \text{AbstrLevel } i$
shows $sA92 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-A9-A93*:
assumes $sA9 \in \text{AbstrLevel } i$
shows $sA93 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S1-A12*:
assumes $sS1 \in \text{AbstrLevel } i$
shows $sA12 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S2-A11*:
assumes $sS2 \in \text{AbstrLevel } i$
shows $sA11 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S3-A21*:
assumes $sS3 \in \text{AbstrLevel } i$
shows $sA21 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S4-A23*:
assumes $sS4 \in \text{AbstrLevel } i$
shows $sA23 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S5-A32*:
assumes $sS5 \in \text{AbstrLevel } i$
shows $sA32 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S6-A22*:
assumes $sS6 \in \text{AbstrLevel } i$
shows $sA22 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S6-A31*:
assumes $sS6 \in \text{AbstrLevel } i$
shows $sA31 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S6-A41*:
assumes $sS6 \in \text{AbstrLevel } i$
shows $sA41 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S7-A42*:
assumes $sS7 \in \text{AbstrLevel } i$
shows $sA42 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S8-A5*:
assumes $sS8 \in \text{AbstrLevel } i$
shows $sA5 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S9-A6*:
assumes $sS9 \in \text{AbstrLevel } i$
shows $sA6 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S10-A71*:
assumes $sS10 \in \text{AbstrLevel } i$
shows $sA71 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S11-A72*:
assumes $sS11 \in \text{AbstrLevel } i$
shows $sA72 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S12-A81*:
assumes $sS12 \in \text{AbstrLevel } i$
shows $sA81 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S12-A91*:
assumes $sS12 \in \text{AbstrLevel } i$
shows $sA91 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S13-A92*:
assumes $sS13 \in \text{AbstrLevel } i$
shows $sA92 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S14-A82*:
assumes $sS14 \in \text{AbstrLevel } i$
shows $sA82 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S15-A93*:
assumes $sS15 \in \text{AbstrLevel } i$
shows $sA93 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S1opt-A11*:
assumes $sS1opt \in \text{AbstrLevel } i$
shows $sA11 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S1opt-A12*:
assumes $sS1opt \in \text{AbstrLevel } i$
shows $sA12 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S4opt-A23*:
assumes $sS4opt \in \text{AbstrLevel } i$
shows $sA23 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S4opt-A32*:
assumes $sS4opt \in \text{AbstrLevel } i$
shows $sA32 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S4opt-A22*:
assumes $sS4opt \in \text{AbstrLevel } i$
shows $sA22 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S4opt-A31*:
assumes $sS4opt \in \text{AbstrLevel } i$
shows $sA31 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S4opt-A41*:
assumes $sS4opt \in \text{AbstrLevel } i$
shows $sA41 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S7opt-A42*:
assumes $sS7opt \in \text{AbstrLevel } i$
shows $sA42 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S7opt-A5*:

assumes $s7_{opt} \in \text{AbstrLevel } i$
shows $sA5 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S11opt-A72*:
assumes $s11_{opt} \in \text{AbstrLevel } i$
shows $sA72 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S11opt-A82*:
assumes $s11_{opt} \in \text{AbstrLevel } i$
shows $sA82 \notin \text{AbstrLevel } i$

lemma *AbstrLevels-S11opt-A93*:
assumes $s11_{opt} \in \text{AbstrLevel } i$
shows $sA93 \notin \text{AbstrLevel } i$

lemma *correctCompositionDiffLevelsA1*: *correctCompositionDiffLevels sA1*

lemma *correctCompositionDiffLevelsA2*: *correctCompositionDiffLevels sA2*

lemma *correctCompositionDiffLevelsA3*: *correctCompositionDiffLevels sA3*

lemma *correctCompositionDiffLevelsA4*: *correctCompositionDiffLevels sA4*

— lemmas *correctCompositionDiffLevelsX* and corresponding proofs
— are identical for all elementary components, they can be constructed automatically

lemma *correctCompositionDiffLevelsA5*: *correctCompositionDiffLevels sA5*

lemma *correctCompositionDiffLevelsA6*: *correctCompositionDiffLevels sA6*

lemma *correctCompositionDiffLevelsA7*: *correctCompositionDiffLevels sA7*

lemma *correctCompositionDiffLevelsA8*: *correctCompositionDiffLevels sA8*

lemma *correctCompositionDiffLevelsA9*: *correctCompositionDiffLevels sA9*

lemma *correctCompositionDiffLevelsA11*: *correctCompositionDiffLevels sA11*

lemma *correctCompositionDiffLevelsA12*: *correctCompositionDiffLevels sA12*

lemma *correctCompositionDiffLevelsA21*: *correctCompositionDiffLevels sA21*

lemma *correctCompositionDiffLevelsA22*: *correctCompositionDiffLevels sA22*

lemma *correctCompositionDiffLevelsA23*: *correctCompositionDiffLevels sA23*

lemma *correctCompositionDiffLevelsA31*: *correctCompositionDiffLevels sA31*

lemma *correctCompositionDiffLevelsA32*: *correctCompositionDiffLevels sA32*

lemma *correctCompositionDiffLevelsA41*: *correctCompositionDiffLevels sA41*

lemma *correctCompositionDiffLevelsA42*: *correctCompositionDiffLevels sA42*

lemma *correctCompositionDiffLevelsA71*: *correctCompositionDiffLevels sA71*

lemma *correctCompositionDiffLevelsA72*: *correctCompositionDiffLevels sA72*

lemma *correctCompositionDiffLevelsA81*: *correctCompositionDiffLevels sA81*

lemma *correctCompositionDiffLevelsA82*: *correctCompositionDiffLevels sA82*

lemma *correctCompositionDiffLevelsA91*: *correctCompositionDiffLevels sA91*

lemma *correctCompositionDiffLevelsA92*: *correctCompositionDiffLevels sA92*

lemma *correctCompositionDiffLevelsA93*: *correctCompositionDiffLevels sA93*

lemma *correctCompositionDiffLevelsS1*: *correctCompositionDiffLevels sS1*

lemma *correctCompositionDiffLevelsS2*: *correctCompositionDiffLevels sS2*

lemma *correctCompositionDiffLevelsS3: correctCompositionDiffLevels sS3*
lemma *correctCompositionDiffLevelsS4: correctCompositionDiffLevels sS4*
lemma *correctCompositionDiffLevelsS5: correctCompositionDiffLevels sS5*
lemma *correctCompositionDiffLevelsS6: correctCompositionDiffLevels sS6*
lemma *correctCompositionDiffLevelsS7: correctCompositionDiffLevels sS7*
lemma *correctCompositionDiffLevelsS8: correctCompositionDiffLevels sS8*
lemma *correctCompositionDiffLevelsS9: correctCompositionDiffLevels sS9*
lemma *correctCompositionDiffLevelsS10: correctCompositionDiffLevels sS10*
lemma *correctCompositionDiffLevelsS11: correctCompositionDiffLevels sS11*
lemma *correctCompositionDiffLevelsS12: correctCompositionDiffLevels sS12*
lemma *correctCompositionDiffLevelsS13: correctCompositionDiffLevels sS13*
lemma *correctCompositionDiffLevelsS14: correctCompositionDiffLevels sS14*
lemma *correctCompositionDiffLevelsS15: correctCompositionDiffLevels sS15*
lemma *correctCompositionDiffLevelsS1opt: correctCompositionDiffLevels sS1opt*
lemma *correctCompositionDiffLevelsS4opt: correctCompositionDiffLevels sS4opt*
lemma *correctCompositionDiffLevelsS7opt: correctCompositionDiffLevels sS7opt*
lemma *correctCompositionDiffLevelsS11opt: correctCompositionDiffLevels sS11opt*
lemma *correctCompositionDiffLevelsSYSTEM-holds:*
correctCompositionDiffLevelsSYSTEM
lemma *correctCompositionVARSYSTEM-holds:*
correctCompositionVARSYSTEM
by (*simp add: correctCompositionVARSYSTEM-def, clarify, case-tac S, (simp add: correctCompositionVAR-def)+*)

lemma *correctDeCompositionVARSYSTEM-holds:*
correctDeCompositionVARSYSTEM
by (*simp add: correctDeCompositionVARSYSTEM-def, clarify, case-tac S, (simp add: correctDeCompositionVAR-def)+*)

4.2 Correct specification of the relations between channels

lemma *OUTfromChCorrect-data1: OUTfromChCorrect data1*
by (*simp add: OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data2: OUTfromChCorrect data2*
by (*metis IN.simps(27) OUT.simps(27) OUTfromCh.simps(2) OUTfromChCorrect-def insertI1*)

lemma *OUTfromChCorrect-data3: OUTfromChCorrect data3*
by (*metis OUTfromCh.simps(3) OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data4: OUTfromChCorrect data4*
by (*metis IN.simps(2) OUT.simps(2) OUTfromCh.simps(4) OUTfromChCorrect-def insertI1 singleton-iff*)

lemma *OUTfromChCorrect-data5: OUTfromChCorrect data5*
by (*simp add: OUTfromChCorrect-def, metis IN.simps(14) OUT.simps(14) insertI1*)

lemma *OUTfromChCorrect-data6: OUTfromChCorrect data6*
by (*simp add: OUTfromChCorrect-def, metis IN.simps(15) OUT.simps(15) insertI1*)

lemma *OUTfromChCorrect-data7: OUTfromChCorrect data7*
by (*simp add: OUTfromChCorrect-def, metis IN.simps(16) OUT.simps(16) insertI1*)

lemma *OUTfromChCorrect-data8: OUTfromChCorrect data8*
by (*simp add: OUTfromChCorrect-def, metis IN.simps(18) OUT.simps(18) insertI1*)

lemma *OUTfromChCorrect-data9: OUTfromChCorrect data9*
by (*simp add: OUTfromChCorrect-def, metis IN.simps(33) OUT.simps(33) singleton-iff*)

lemma *OUTfromChCorrect-data10: OUTfromChCorrect data10*
by (*simp add: OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data11: OUTfromChCorrect data11*
by (*simp add: OUTfromChCorrect-def, metis (full-types) IN.simps(2) OUT.simps(2) OUT.simps(31) Un-empty-right Un-insert-left Un-insert-right insertI1*)

lemma *OUTfromChCorrect-data12: OUTfromChCorrect data12*
by (*simp add: OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data13: OUTfromChCorrect data13*
by (*simp add: OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data14: OUTfromChCorrect data14*
by (*metis OUTfromCh.simps(14) OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data15: OUTfromChCorrect data15*
by (*metis OUTfromCh.simps(15) OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data16: OUTfromChCorrect data16*
by (*metis OUTfromCh.simps(16) OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data17: OUTfromChCorrect data17*

proof –

have *data17 ∈ OUT sA71 ∧ data15 ∈ IN sA71*

by (*metis IN.simps(19) OUT.simps(19) insertI1*)

thus *?thesis* **by** (*metis IN.simps(19) OUTfromCh.simps(17) OUTfromChCorrect-def*)

qed

lemma *OUTfromChCorrect-data18: OUTfromChCorrect data18*

by (*simp add: OUTfromChCorrect-def, metis IN.simps(20) OUT.simps(20) insertI1*)

lemma *OUTfromChCorrect-data19: OUTfromChCorrect data19*

by (*metis OUTfromCh.simps(19) OUTfromChCorrect-def*)

lemma *OUTfromChCorrect-data20: OUTfromChCorrect data20*

by (*simp add: OUTfromChCorrect-def, metis IN.simps(21) OUT.simps(21) insertI1 insert-subset subset-insertI*)

lemma *OUTfromChCorrect-data21: OUTfromChCorrect data21*

by (*simp add: OUTfromChCorrect-def, metis (full-types)*

IN.simps(22) OUT.simps(22) insertI1 insert-subset subset-insertI)

lemma *OUTfromChCorrect-data22: OUTfromChCorrect data22*

by (*simp add: OUTfromChCorrect-def, metis (full-types) IN.simps(23) OUT.simps(23) insertI1*)

lemma *OUTfromChCorrect-data23: OUTfromChCorrect data23*

by (*simp add: OUTfromChCorrect-def, metis (full-types) IN.simps(9) OUT.simps(9) insert-subset subset-insertI*)

lemma *OUTfromChCorrect-data24: OUTfromChCorrect data24*

by (*simp add: OUTfromChCorrect-def, metis IN.simps(9) OUT.simps(9) insertI1 insert-subset subset-insertI*)

lemma *OUTfromChCorrectSYSTEM-holds: OUTfromChCorrectSYSTEM*

by (*simp add: OUTfromChCorrectSYSTEM-def, clarify, case-tac x,*

simp add: OUTfromChCorrect-data1, simp add: OUTfromChCorrect-data2,

simp add: OUTfromChCorrect-data3, simp add: OUTfromChCorrect-data4,

simp add: OUTfromChCorrect-data5, simp add: OUTfromChCorrect-data6,

simp add: OUTfromChCorrect-data7, simp add: OUTfromChCorrect-data8,

simp add: OUTfromChCorrect-data9, simp add: OUTfromChCorrect-data10,

simp add: OUTfromChCorrect-data11, simp add: OUTfromChCorrect-data12,

simp add: OUTfromChCorrect-data13, simp add: OUTfromChCorrect-data14,

simp add: OUTfromChCorrect-data15, simp add: OUTfromChCorrect-data16,

simp add: OUTfromChCorrect-data17, simp add: OUTfromChCorrect-data18,

simp add: OUTfromChCorrect-data19, simp add: OUTfromChCorrect-data20,

simp add: OUTfromChCorrect-data21, simp add: OUTfromChCorrect-data22,

simp add: OUTfromChCorrect-data23, simp add: OUTfromChCorrect-data24)

lemma *OUTfromVCorrect1-data1: OUTfromVCorrect1 data1*

by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data2: OUTfromVCorrect1 data2*

by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data3: OUTfromVCorrect1 data3*

proof –

have $data3 \in OUT\ sA41 \wedge stA4 \in VAR\ sA41$
by (*metis* $OUT.simps(17)$ $VAR.simps(17)$ *insertI1*)
thus *?thesis* **by** (*metis* $OUTfromV.simps(3)$ $OUTfromVCorrect1-def$ $VAR.simps(17)$)

qed

lemma $OUTfromVCorrect1-data4: OUTfromVCorrect1\ data4$
by (*simp* *add: OUTfromVCorrect1-def, metis (full-types) OUT.simps(2) VAR.simps(2) insertI1*)

lemma $OUTfromVCorrect1-data5: OUTfromVCorrect1\ data5$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data6: OUTfromVCorrect1\ data6$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data7: OUTfromVCorrect1\ data7$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data8: OUTfromVCorrect1\ data8$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data9: OUTfromVCorrect1\ data9$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data10: OUTfromVCorrect1\ data10$
proof –

have $data10 \in OUT\ sA12 \wedge stA1 \in VAR\ sA12$
by (*metis* $OUT.simps(11)$ $VAR.simps(11)$ *insertI1*)
thus *?thesis* **by** (*metis* $OUT.simps(26)$ $OUTfromV.simps(10)$ $OUTfromVCorrect1-def$ $VAR.simps(26)$ *insertI1*)

qed

lemma $OUTfromVCorrect1-data11: OUTfromVCorrect1\ data11$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data12: OUTfromVCorrect1\ data12$
proof –

have $data12 \in OUT\ sA22 \wedge stA2 \in VAR\ sA22$
by (*metis* (*full-types*) $OUT.simps(13)$ $VAR.simps(13)$ *insertCI*)
thus *?thesis* **by** (*metis* $OUTfromV.simps(12)$ $OUTfromVCorrect1-def$ $VAR.simps(13)$)

qed

lemma $OUTfromVCorrect1-data13: OUTfromVCorrect1\ data13$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma $OUTfromVCorrect1-data14: OUTfromVCorrect1\ data14$
by (*simp* *add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data15: OUTfromVCorrect1 data15*
proof –
 have *A6ch:data15 ∈ OUT sA6 ∧ stA6 ∈ VAR sA6*
 by (*metis OUT.simps(6) VAR.simps(6) insertI1*)
 thus *?thesis* **by** (*simp add: OUTfromVCorrect1-def, metis A6ch*)
qed

lemma *OUTfromVCorrect1-data16: OUTfromVCorrect1 data16*
proof –
 have *A6ch:data16 ∈ OUT sA6 ∧ stA6 ∈ VAR sA6*
 by (*metis (full-types) OUT.simps(6) VAR.simps(6) insertCI*)
 thus *?thesis* **by** (*simp add: OUTfromVCorrect1-def, metis A6ch*)
qed

lemma *OUTfromVCorrect1-data17: OUTfromVCorrect1 data17*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data18: OUTfromVCorrect1 data18*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data19: OUTfromVCorrect1 data19*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data20: OUTfromVCorrect1 data20*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data21: OUTfromVCorrect1 data21*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data22: OUTfromVCorrect1 data22*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data23: OUTfromVCorrect1 data23*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1-data24: OUTfromVCorrect1 data24*
by (*simp add: OUTfromVCorrect1-def*)

lemma *OUTfromVCorrect1SYSTEM-holds: OUTfromVCorrect1SYSTEM*
by (*simp add: OUTfromVCorrect1SYSTEM-def, clarify, case-tac x,*
simp add: OUTfromVCorrect1-data1, simp add: OUTfromVCorrect1-data2,
simp add: OUTfromVCorrect1-data3, simp add: OUTfromVCorrect1-data4,
simp add: OUTfromVCorrect1-data5, simp add: OUTfromVCorrect1-data6,
simp add: OUTfromVCorrect1-data7, simp add: OUTfromVCorrect1-data8,
simp add: OUTfromVCorrect1-data9, simp add: OUTfromVCorrect1-data10,
simp add: OUTfromVCorrect1-data11, simp add: OUTfromVCorrect1-data12,
simp add: OUTfromVCorrect1-data13, simp add: OUTfromVCorrect1-data14,
simp add: OUTfromVCorrect1-data15, simp add: OUTfromVCorrect1-data16,

simp add: OUTfromVCorrect1-data17, simp add: OUTfromVCorrect1-data18,
simp add: OUTfromVCorrect1-data19, simp add: OUTfromVCorrect1-data20,
simp add: OUTfromVCorrect1-data21, simp add: OUTfromVCorrect1-data22,
simp add: OUTfromVCorrect1-data23, simp add: OUTfromVCorrect1-data24)

lemma *OUTfromVCorrect2SYSTEM: OUTfromVCorrect2SYSTEM*
by (*simp add: OUTfromVCorrect2SYSTEM-def, auto, case-tac x,*
((simp add: OUTfromVCorrect2-def, auto, case-tac v, auto) |
(simp add: OUTfromVCorrect2-def))+)

lemma *OUTfromV-VARto-holds:*
OUTfromV-VARto
by (*simp add: OUTfromV-VARto-def, auto, (case-tac x, auto), (case-tac v, auto)*)

lemma *VARfromCorrectSYSTEM-holds:*
VARfromCorrectSYSTEM
by (*simp add: VARfromCorrectSYSTEM-def AbstrLevel0 AbstrLevel1*)

lemma *VARtoCorrectSYSTEM-holds:*
VARtoCorrectSYSTEM
by (*simp add: VARtoCorrectSYSTEM-def AbstrLevel0 AbstrLevel1*)

lemma *VARusefulSYSTEM-holds:*
VARusefulSYSTEM
by (*simp add: VARusefulSYSTEM-def, auto, case-tac v, auto*)

4.3 Elementary components

— On the abstraction level 0 only the components sA5 and sA6 are elementary

lemma *NOT-elementaryCompDD-sA1: \neg elementaryCompDD sA1*

proof —

have *outSetCorelated data2 \cap outSetCorelated data10 = {}*
by (*metis OUTfromV.simps(2) inf-bot-left outSetCorelatedEmpty1*)
thus *?thesis by (simp add: elementaryCompDD-def)*

qed

lemma *NOT-elementaryCompDD-sA2: \neg elementaryCompDD sA2*

proof —

have *outSetCorelated data5 \cap outSetCorelated data11 = {}*
by (*metis OUTfromV.simps(5) inf-bot-right inf-commute outSetCorelatedEmpty1*)
thus *?thesis by (simp add: elementaryCompDD-def)*

qed

lemma *NOT-elementaryCompDD-sA3: \neg elementaryCompDD sA3*

proof —

have *outSetCorelated data6 \cap outSetCorelated data7 = {}*
by (*metis OUTfromV.simps(7) inf-bot-right outSetCorelatedEmpty1*)
thus *?thesis by (simp add: elementaryCompDD-def)*

qed

lemma *NOT-elementaryCompDD-sA4*: \neg elementaryCompDD sA4

proof –

have outSetCorelated data3 \cap outSetCorelated data8 = {}

by (metis OUTfromV.simps(8) inf-bot-left inf-commute outSetCorelatedEmpty1)

thus ?thesis **by** (simp add: elementaryCompDD-def)

qed

lemma *elementaryCompDD-sA5*: elementaryCompDD sA5

by (simp add: elementaryCompDD-def)

lemma *elementaryCompDD-sA6*: elementaryCompDD sA6

proof –

have oSet15:outSetCorelated data15 \neq {}

by (simp add: outSetCorelated-def, auto)

have oSet16:outSetCorelated data16 \neq {}

by (simp add: outSetCorelated-def, auto)

have outSetCorelated data15 \cap outSetCorelated data16 \neq {}

by (simp add: outSetCorelated-def, auto)

with oSet15 oSet16 **show** ?thesis **by** (simp add: elementaryCompDD-def, auto)

qed

lemma *NOT-elementaryCompDD-sA7*: \neg elementaryCompDD sA7

proof –

have outSetCorelated data17 \cap outSetCorelated data18 = {}

by (metis (full-types) OUTfromV.simps(17) disjoint-iff-not-equal empty-iff outSetCorelatedEmpty1)

thus ?thesis **by** (simp add: elementaryCompDD-def)

qed

lemma *NOT-elementaryCompDD-sA8*: \neg elementaryCompDD sA8

proof –

have outSetCorelated data20 \cap outSetCorelated data21 = {}

by (metis OUTfromV.simps(21) inf-bot-right outSetCorelatedEmpty1)

thus ?thesis **by** (simp add: elementaryCompDD-def)

qed

lemma *NOT-elementaryCompDD-sA9*: \neg elementaryCompDD sA9

proof –

have outSetCorelated data23 \cap outSetCorelated data24 = {}

by (metis (full-types) OUTfromV.simps(23) disjoint-iff-not-equal empty-iff outSetCorelatedEmpty1)

thus ?thesis **by** (simp add: elementaryCompDD-def)

qed

— On the abstraction level 1 all components are elementary

lemma *elementaryCompDD-sA11: elementaryCompDD sA11*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA12: elementaryCompDD sA12*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA21: elementaryCompDD sA21*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA22: elementaryCompDD sA22*
proof –

have *oSet4:outSetCorelated data4 \neq {}*
by (*simp add: outSetCorelated-def, auto*)
have *oSet12:outSetCorelated data12 \neq {}*
by (*simp add: outSetCorelated-def, auto*)
have *outSetCorelated data4 \cap outSetCorelated data12 \neq {}*
by (*simp add: outSetCorelated-def, auto*)
with *oSet4 oSet12* **show** *?thesis*
by (*simp add: elementaryCompDD-def, auto*)

qed

lemma *elementaryCompDD-sA23: elementaryCompDD sA23*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA31: elementaryCompDD sA31*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA32: elementaryCompDD sA32*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA41: elementaryCompDD sA41*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA42: elementaryCompDD sA42*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA71: elementaryCompDD sA71*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA72: elementaryCompDD sA72*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA81: elementaryCompDD sA81*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA82: elementaryCompDD sA82*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA91: elementaryCompDD sA91*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA92: elementaryCompDD sA92*
by (*simp add: elementaryCompDD-def*)

lemma *elementaryCompDD-sA93: elementaryCompDD sA93*
by (*simp add: elementaryCompDD-def*)

4.4 Source components

— Abstraction level 0

lemma *A5-NotDSource-level0: isNotDSource level0 sA5*
by (*simp add: isNotDSource-def, auto, case-tac Z, auto*)

lemma *DSourcesA1-L0: DSources level0 sA1 = {}*
by (*simp add: DSources-def, auto, case-tac x, auto*)

lemma *DSourcesA2-L0: DSources level0 sA2 = { sA1, sA4 }*
by (*simp add: DSources-def AbstrLevel0, auto*)

lemma *DSourcesA3-L0: DSources level0 sA3 = { sA2 }*
by (*simp add: DSources-def AbstrLevel0, auto*)

lemma *DSourcesA4-L0: DSources level0 sA4 = { sA3 }*
by (*simp add: DSources-def AbstrLevel0, auto*)

lemma *DSourcesA5-L0: DSources level0 sA5 = { sA4 }*
by (*simp add: DSources-def AbstrLevel0, auto*)

lemma *DSourcesA6-L0: DSources level0 sA6 = {}*
by (*simp add: DSources-def, auto, case-tac x, auto*)

lemma *DSourcesA7-L0: DSources level0 sA7 = { sA6 }*
by (*simp add: DSources-def AbstrLevel0, auto*)

lemma *DSourcesA8-L0: DSources level0 sA8 = { sA7, sA9 }*
by (*simp add: DSources-def AbstrLevel0, force*)

lemma *DSourcesA9-L0: DSources level0 sA9 = { sA8 }*
by (*simp add: DSources-def AbstrLevel0, auto*)

lemma *A1-DAcc-level0: DAcc level0 sA1 = { sA2 }*
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A2-DAcc-level0: DAcc level0 sA2 = { sA3 }*
by (*simp add: DAcc-def AbstrLevel0, force*)

lemma *A3-DAcc-level0*: $D\text{Acc level0 } sA3 = \{ sA4 \}$
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A4-DAcc-level0*: $D\text{Acc level0 } sA4 = \{ sA2, sA5 \}$
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A5-DAcc-level0*: $D\text{Acc level0 } sA5 = \{ \}$
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A6-DAcc-level0*: $D\text{Acc level0 } sA6 = \{ sA7 \}$
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A7-DAcc-level0*: $D\text{Acc level0 } sA7 = \{ sA8 \}$
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A8-DAcc-level0*: $D\text{Acc level0 } sA8 = \{ sA9 \}$
by (*simp add: DAcc-def AbstrLevel0, auto*)

lemma *A9-DAcc-level0*: $D\text{Acc level0 } sA9 = \{ sA8 \}$
by (*simp add: DAcc-def AbstrLevel0, force*)

lemma *A8-NSources*:
 $\forall C \in (\text{AbstrLevel level0}). (C \neq sA9 \wedge C \neq sA8 \longrightarrow sA8 \notin (\text{Sources level0 } C))$
by (*metis A8-DAcc-level0 A9-DAcc-level0 singleDSourceLoop*)

lemma *A9-NSources*:
 $\forall C \in (\text{AbstrLevel level0}). (C \neq sA9 \wedge C \neq sA8 \longrightarrow sA9 \notin (\text{Sources level0 } C))$
by (*metis A8-DAcc-level0 A9-DAcc-level0 singleDSourceLoop*)

lemma *A7-Acc*:
 $(\text{Acc level0 } sA7) = \{ sA8, sA9 \}$
by (*metis A7-DAcc-level0 A8-DAcc-level0 A9-DAcc-level0 AccDef AccSigleLoop insert-commute*)

lemma *A7-NSources*:
 $\forall C \in (\text{AbstrLevel level0}). (C \neq sA9 \wedge C \neq sA8 \longrightarrow sA7 \notin (\text{Sources level0 } C))$
by (*metis A7-Acc Acc-Sources insert-iff singleton-iff*)

lemma *A5-Acc*: $(\text{Acc level0 } sA5) = \{ \}$
by (*metis A5-NotDSource-level0 isNotDSource-EmptyAcc*)

lemma *A6-Acc*:
 $(\text{Acc level0 } sA6) = \{ sA7, sA8, sA9 \}$

proof –

have *daA6*: $D\text{Acc level0 } sA6 = \{ sA7 \}$ **by** (*rule A6-DAcc-level0*)

hence $(\bigcup S \in (D\text{Acc level0 } sA6). (\text{Acc level0 } S)) = (\text{Acc level0 } sA7)$ **by** *simp*

hence *aA6*: $(\bigcup S \in (D\text{Acc level0 } sA6). (\text{Acc level0 } S)) = \{ sA8, sA9 \}$ **by** (*simp add: A7-Acc*)

have $(\text{Acc level0 } sA6) = (D\text{Acc level0 } sA6) \cup (\bigcup S \in (D\text{Acc level0 } sA6). (\text{Acc$

level0 S))
 by (rule AccDef)
 with daA6 aA6 show ?thesis by auto
 qed

lemma A6-NSources:
 $\forall C \in (\text{AbstrLevel level0}). (C \neq sA9 \wedge C \neq sA8 \wedge C \neq sA7 \longrightarrow sA6 \notin (\text{Sources level0 } C))$
 by (metis (full-types) A6-Acc A7-Acc Acc-SourcesNOT insert-iff singleton-iff)

lemma SourcesA1-L0: Sources level0 sA1 = {}
 by (simp add: DSourcesA1-L0 DSourcesEmptySources)

lemma SourcesA2-L0: Sources level0 sA2 = { sA1, sA2, sA3, sA4 }

proof

show Sources level0 sA2 \subseteq {sA1, sA2, sA3, sA4}

proof –

have A2level0:sA2 \in (AbstrLevel level0) by (simp add: AbstrLevel0)

have sgA5:sA5 \notin Sources level0 sA2

by (metis A5-NotDSource-level0 DSource-level NoDSourceNoSource
 allNotDSource-NotSource isNotSource-Sources)

from A2level0 have sgA6:sA6 \notin Sources level0 sA2 by (simp add: A6-NSources)

from A2level0 have sgA7:sA7 \notin Sources level0 sA2 by (simp add: A7-NSources)

from A2level0 have sgA8:sA8 \notin Sources level0 sA2 by (simp add: A8-NSources)

from A2level0 have sgA9:sA9 \notin Sources level0 sA2 by (simp add: A9-NSources)

have Sources level0 sA2 \subseteq {sA1, sA2, sA3, sA4, sA5, sA6, sA7, sA8, sA9}

by (metis AbstrLevel0 SourcesLevelX)

with sgA5 sgA6 sgA7 sgA8 sgA9 show Sources level0 sA2 \subseteq {sA1, sA2, sA3,
 sA4}

by blast

qed

next

show {sA1, sA2, sA3, sA4} \subseteq Sources level0 sA2

proof –

have dsA4:{ sA3 } \subseteq Sources level0 sA2

by (metis DSource-Sources DSourcesA2-L0 DSourcesA4-L0
 Sources-DSources insertI1 insert-commute subset-trans)

have { sA2 } \subseteq Sources level0 sA2

by (metis DSource-Sources DSourcesA2-L0 DSourcesA3-L0
 DSourcesA4-L0 Sources-DSources insertI1
 insert-commute subset-trans)

with dsA4 show {sA1, sA2, sA3, sA4} \subseteq Sources level0 sA2

by (metis DSourcesA2-L0 Sources-DSources insert-subset)

qed

qed

lemma SourcesA3-L0: Sources level0 sA3 = { sA1, sA2, sA3, sA4 }

proof

show Sources level0 sA3 \subseteq {sA1, sA2, sA3, sA4}

proof –
have $a2: \text{Sources level0 } sA2 = \{sA1, sA2, sA3, sA4\}$ **by** (*simp add: SourcesA2-L0*)
have $\{sA2\} \subseteq \text{DSources level0 } sA3$ **by** (*simp add: DSourcesA3-L0*)
with $a2$ **show** $\text{Sources level0 } sA3 \subseteq \{sA1, sA2, sA3, sA4\}$
by (*metis DSource-Sources DSourcesA2-L0 DSourcesA4-L0 insertI1 insert-commute subset-trans*)
qed
next
show $\{sA1, sA2, sA3, sA4\} \subseteq \text{Sources level0 } sA3$
by (*metis (full-types) DSource-Sources DSourcesA3-L0 SourcesA2-L0 insertI1*)
qed

lemma *SourcesA4-L0*: $\text{Sources level0 } sA4 = \{sA1, sA2, sA3, sA4\}$

proof –
have $A3s: \text{Sources level0 } sA3 = \{sA1, sA2, sA3, sA4\}$ **by** (*rule SourcesA3-L0*)
have $\text{Sources level0 } sA4 = \{sA3\} \cup \text{Sources level0 } sA3$
by (*metis DSourcesA4-L0 Sources-singleDSource*)
with $A3s$ **show** *?thesis* **by** *auto*
qed

lemma *SourcesA5-L0*: $\text{Sources level0 } sA5 = \{sA1, sA2, sA3, sA4\}$

proof –
have $A4s: \text{Sources level0 } sA4 = \{sA1, sA2, sA3, sA4\}$ **by** (*rule SourcesA4-L0*)
have $\text{Sources level0 } sA5 = \{sA4\} \cup \text{Sources level0 } sA4$
by (*metis DSourcesA5-L0 Sources-singleDSource*)
with $A4s$ **show** *?thesis* **by** *auto*
qed

lemma *SourcesA6-L0*: $\text{Sources level0 } sA6 = \{\}$

by (*simp add: DSourcesA6-L0 DSourcesEmptySources*)

lemma *SourcesA7-L0*: $\text{Sources level0 } sA7 = \{sA6\}$

by (*metis DSourcesA7-L0 SourcesA6-L0 SourcesEmptyDSources SourcesOnlyD-Sources singleton-iff*)

lemma *SourcesA8-L0*: $\text{Sources level0 } sA8 = \{sA6, sA7, sA8, sA9\}$

proof –
have $dA8: \text{DSources level0 } sA8 = \{sA7, sA9\}$ **by** (*rule DSourcesA8-L0*)
have $dA9: \text{DSources level0 } sA9 = \{sA8\}$ **by** (*rule DSourcesA9-L0*)
have $(\text{Sources level0 } sA8) = (\text{DSources level0 } sA8) \cup (\bigcup S \in (\text{DSources level0 } sA8). (\text{Sources level0 } S))$
by (*rule SourcesDef*)
hence $\text{sourcesA8}:(\text{Sources level0 } sA8) = (\{sA7, sA9, sA6\} \cup (\text{Sources level0 } sA9))$
by (*simp add: DSourcesA8-L0 SourcesA7-L0, auto*)
have $(\text{Sources level0 } sA9) = (\text{DSources level0 } sA9) \cup (\bigcup S \in (\text{DSources level0 } sA9). (\text{Sources level0 } S))$
by (*rule SourcesDef*)

hence (*Sources level0 sA9*) = ($\{sA8\} \cup$ (*Sources level0 sA8*))
by (*simp add: DSourcesA9-L0*)
with sourcesA8 have (*Sources level0 sA8*) = $\{sA7, sA9, sA6\} \cup \{sA8\} \cup \{sA8,$
sA9 $\}$
by (*metis SourcesLoop*)
thus *?thesis* **by** *auto*
qed

lemma *SourcesA9-L0: Sources level0 sA9 = { sA6, sA7, sA8, sA9 }*
proof –
have (*Sources level0 sA9*) = (*DSources level0 sA9*) \cup ($\bigcup S \in$ (*DSources level0*
sA9). (*Sources level0 S*))
by (*rule SourcesDef*)
hence *sourcesA9:(Sources level0 sA9)* = ($\{sA8\} \cup$ (*Sources level0 sA8*))
by (*simp add: DSourcesA9-L0*)
thus *?thesis* **by** (*metis SourcesA8-L0 Un-insert-right insert-absorb2 insert-is-Un*)

qed

— Abstraction level 1

lemma *A12-NotSource-level1: isNotDSource level1 sA12*
by (*simp add: isNotDSource-def, auto, case-tac Z, auto*)

lemma *A21-NotSource-level1: isNotDSource level1 sA21*
by (*simp add: isNotDSource-def, auto, case-tac Z, auto*)

lemma *A5-NotSource-level1: isNotDSource level1 sA5*
by (*simp add: isNotDSource-def, auto, case-tac Z, auto*)

lemma *A92-NotSource-level1: isNotDSource level1 sA92*
by (*simp add: isNotDSource-def, auto, case-tac Z, auto*)

lemma *A93-NotSource-level1: isNotDSource level1 sA93*
by (*simp add: isNotDSource-def, auto, case-tac Z, auto*)

lemma *A11-DAcc-level1: DAcc level1 sA11 = { sA21, sA22, sA23 }*
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A12-DAcc-level1: DAcc level1 sA12 = {}*
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A21-DAcc-level1: DAcc level1 sA21 = {}*
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A22-DAcc-level1: DAcc level1 sA22 = {sA31}*
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A23-DAcc-level1*: $D\text{Acc level1 } sA23 = \{sA32\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A31-DAcc-level1*: $D\text{Acc level1 } sA31 = \{sA41\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A32-DAcc-level1*: $D\text{Acc level1 } sA32 = \{sA41\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A41-DAcc-level1*: $D\text{Acc level1 } sA41 = \{sA22\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A42-DAcc-level1*: $D\text{Acc level1 } sA42 = \{sA5\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A5-DAcc-level1*: $D\text{Acc level1 } sA5 = \{\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A6-DAcc-level1*: $D\text{Acc level1 } sA6 = \{sA71, sA72\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A71-DAcc-level1*: $D\text{Acc level1 } sA71 = \{sA81\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A72-DAcc-level1*: $D\text{Acc level1 } sA72 = \{sA82\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A81-DAcc-level1*: $D\text{Acc level1 } sA81 = \{sA91, sA92\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A82-DAcc-level1*: $D\text{Acc level1 } sA82 = \{sA93\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A91-DAcc-level1*: $D\text{Acc level1 } sA91 = \{sA81\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A92-DAcc-level1*: $D\text{Acc level1 } sA92 = \{\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A93-DAcc-level1*: $D\text{Acc level1 } sA93 = \{\}$
by (*simp add: DAcc-def AbstrLevel1, auto*)

lemma *A42-NSources-L1*:
 $\forall C \in (\text{AbstrLevel level1}). C \neq sA5 \longrightarrow sA42 \notin (\text{Sources level1 } C)$
by (*metis A42-DAcc-level1 A5-NotSource-level1 singleDSourceEmpty4isNotSource*)

lemma *A5-NotSourceSet-level1* :
 $\forall C \in (\text{AbstrLevel level1}). sA5 \notin (\text{Sources level1 } C)$

by (*metis A5-NotSource-level1 isNotSource-Sources*)

lemma *A92-NotSourceSet-level1* :

$\forall C \in (\text{AbstrLevel level1}). sA92 \notin (\text{Sources level1 } C)$

by (*metis A92-NotSource-level1 isNotSource-Sources*)

lemma *A93-NotSourceSet-level1* :

$\forall C \in (\text{AbstrLevel level1}). sA93 \notin (\text{Sources level1 } C)$

by (*metis A93-NotSource-level1 isNotSource-Sources*)

lemma *DSourcesA11-L1*: $DSources\ level1\ sA11 = \{\}$

by (*simp add: DSources-def, auto, case-tac x, auto*)

lemma *DSourcesA12-L1*: $DSources\ level1\ sA12 = \{\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA21-L1*: $DSources\ level1\ sA21 = \{sA11\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA22-L1*: $DSources\ level1\ sA22 = \{sA11, sA41\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA23-L1*: $DSources\ level1\ sA23 = \{sA11\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA31-L1*: $DSources\ level1\ sA31 = \{sA22\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA32-L1*: $DSources\ level1\ sA32 = \{sA23\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA41-L1*: $DSources\ level1\ sA41 = \{sA31, sA32\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA42-L1*: $DSources\ level1\ sA42 = \{\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA5-L1*: $DSources\ level1\ sA5 = \{sA42\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA6-L1*: $DSources\ level1\ sA6 = \{\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA71-L1*: $DSources\ level1\ sA71 = \{sA6\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA72-L1*: $DSources\ level1\ sA72 = \{sA6\}$

by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA81-L1*: $DSources\ level1\ sA81 = \{ sA71, sA91 \}$
by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA82-L1*: $DSources\ level1\ sA82 = \{ sA72 \}$
by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA91-L1*: $DSources\ level1\ sA91 = \{ sA81 \}$
by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA92-L1*: $DSources\ level1\ sA92 = \{ sA81 \}$
by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *DSourcesA93-L1*: $DSources\ level1\ sA93 = \{ sA82 \}$
by (*simp add: DSources-def AbstrLevel1, auto*)

lemma *A82-Acc*: $(Acc\ level1\ sA82) = \{sA93\}$
by (*metis A82-DAcc-level1 A93-NotSource-level1 singleDSourceEmpty-Acc*)

lemma *A82-NSources-L1*:
 $\forall C \in (AbstrLevel\ level1). (C \neq sA93 \longrightarrow sA82 \notin (Sources\ level1\ C))$
by (*metis A82-Acc Acc-Sources singleton-iff*)

lemma *A72-Acc*: $(Acc\ level1\ sA72) = \{sA82, sA93\}$

proof –

have *daA72*: $DAcc\ level1\ sA72 = \{ sA82 \}$ **by** (*rule A72-DAcc-level1*)
hence $(\bigcup S \in (DAcc\ level1\ sA72). (Acc\ level1\ S)) = (Acc\ level1\ sA82)$ **by** *simp*
hence *aA72*: $(\bigcup S \in (DAcc\ level1\ sA72). (Acc\ level1\ S)) = \{ sA93 \}$ **by** (*simp add: A82-Acc*)

have $(Acc\ level1\ sA72) = (DAcc\ level1\ sA72) \cup (\bigcup S \in (DAcc\ level1\ sA72). (Acc\ level1\ S))$

by (*rule AccDef*)

with *daA72 aA72* **show** *?thesis* **by** *auto*

qed

lemma *A72-NSources-L1*:
 $\forall C \in (AbstrLevel\ level1). (C \neq sA93 \wedge C \neq sA82 \longrightarrow sA72 \notin (Sources\ level1\ C))$
by (*metis A72-Acc Acc-Sources insert-iff singleton-iff*)

lemma *A92-Acc*: $(Acc\ level1\ sA92) = \{\}$
by (*metis A92-NotSource-level1 isNotDSource-EmptyAcc*)

lemma *A92-NSources-L1*:
 $\forall C \in (AbstrLevel\ level1). (sA92 \notin (Sources\ level1\ C))$
by (*metis A92-NotSourceSet-level1*)

lemma *A91-Acc*: $(Acc\ level1\ sA91) = \{sA81, sA91, sA92\}$

proof –

have $da91$: $D\text{Acc level1 } sA91 = \{ sA81 \}$ **by** (rule $A91\text{-}D\text{Acc-level1}$)
hence $a91$: $(\bigcup S \in (D\text{Acc level1 } sA91). (\text{Acc level1 } S)) = (\text{Acc level1 } sA81)$ **by**
simp
have $(\text{Acc level1 } sA91) = (D\text{Acc level1 } sA91) \cup (\bigcup S \in (D\text{Acc level1 } sA91). (\text{Acc level1 } S))$ **by** (rule AccDef)
with $da91$ $a91$ **have** $acc91$: $(\text{Acc level1 } sA91) = \{ sA81 \} \cup (\text{Acc level1 } sA81)$ **by**
simp
have $da81$: $D\text{Acc level1 } sA81 = \{ sA91, sA92 \}$ **by** (rule $A81\text{-}D\text{Acc-level1}$)
hence $a81$: $(\bigcup S \in (D\text{Acc level1 } sA81). (\text{Acc level1 } S)) = (\text{Acc level1 } sA92) \cup (\text{Acc level1 } sA91)$ **by** *auto*
have $(\text{Acc level1 } sA81) = (D\text{Acc level1 } sA81) \cup (\bigcup S \in (D\text{Acc level1 } sA81). (\text{Acc level1 } S))$ **by** (rule AccDef)
with $da81$ $a81$ **have** $acc81$: $(\text{Acc level1 } sA81) = \{ sA91, sA92 \} \cup (\text{Acc level1 } sA91)$
by (metis $A92\text{-}Acc$ *sup-bot.left-neutral*)
from $acc91$ $acc81$ **have** $(\text{Acc level1 } sA91) = \{ sA81 \} \cup \{ sA91, sA92 \} \cup \{ sA91, sA81 \}$
by (metis AccLoop)
thus *?thesis* **by** *auto*
qed

lemma $A91\text{-}N\text{Sources-L1}$:

$\forall C \in (\text{AbstrLevel level1}). (C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA81 \longrightarrow sA91 \notin (\text{Sources level1 } C))$

proof –

have $\forall C \in (\text{AbstrLevel level1}). (C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA81 \longrightarrow (C \notin (\text{Acc level1 } sA91)))$

by (metis $A91\text{-}Acc$ *insert-iff singleton-iff*)

thus *?thesis* **by** (metis Acc-SourcesNOT)

qed

lemma $A81\text{-}Acc$: $(\text{Acc level1 } sA81) = \{ sA81, sA91, sA92 \}$

proof –

have $da91$: $D\text{Acc level1 } sA91 = \{ sA81 \}$ **by** (rule $A91\text{-}D\text{Acc-level1}$)

hence $a91$: $(\bigcup S \in (D\text{Acc level1 } sA91). (\text{Acc level1 } S)) = (\text{Acc level1 } sA81)$ **by**

simp

have $(\text{Acc level1 } sA91) = (D\text{Acc level1 } sA91) \cup (\bigcup S \in (D\text{Acc level1 } sA91). (\text{Acc level1 } S))$ **by** (rule AccDef)

with $da91$ $a91$ **have** $acc91$: $(\text{Acc level1 } sA91) = \{ sA81 \} \cup (\text{Acc level1 } sA81)$ **by**

simp

have $da81$: $D\text{Acc level1 } sA81 = \{ sA91, sA92 \}$ **by** (rule $A81\text{-}D\text{Acc-level1}$)

hence $a81$: $(\bigcup S \in (D\text{Acc level1 } sA81). (\text{Acc level1 } S)) = (\text{Acc level1 } sA92) \cup (\text{Acc level1 } sA91)$ **by** *auto*

have $(\text{Acc level1 } sA81) = (D\text{Acc level1 } sA81) \cup (\bigcup S \in (D\text{Acc level1 } sA81). (\text{Acc level1 } S))$ **by** (rule AccDef)

with $da81$ $a81$ **have** $acc81$: $(\text{Acc level1 } sA81) = \{ sA91, sA92 \} \cup (\text{Acc level1 } sA91)$

by (metis $A92\text{-}Acc$ *sup-bot.left-neutral*)

from $acc81$ $acc91$ **have** $(\text{Acc level1 } sA81) = \{ sA91, sA92 \} \cup \{ sA81 \} \cup$

{sA81, sA91}
by (metis AccLoop)
thus ?thesis **by** auto
qed

lemma A81-NSources-L1:

$\forall C \in (\text{AbstrLevel level1}). (C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA81 \longrightarrow sA81 \notin (\text{Sources level1 } C))$

proof –

have $\forall C \in (\text{AbstrLevel level1}). (C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA81 \longrightarrow (C \notin (\text{Acc level1 } sA81)))$

by (metis A81-Acc insert-iff singleton-iff)

thus ?thesis **by** (metis Acc-SourcesNOT)

qed

lemma A71-Acc: $(\text{Acc level1 } sA71) = \{sA81, sA91, sA92\}$

proof –

have da71: $D\text{Acc level1 } sA71 = \{sA81\}$ **by** (rule A71-DAcc-level1)

hence a71: $(\bigcup S \in (D\text{Acc level1 } sA71). (\text{Acc level1 } S)) = (\text{Acc level1 } sA81)$ **by** simp

have $(\text{Acc level1 } sA71) = (D\text{Acc level1 } sA71) \cup (\bigcup S \in (D\text{Acc level1 } sA71). (\text{Acc level1 } S))$ **by** (rule AccDef)

with da71 a71 **show** ?thesis **by** (metis A91-Acc A91-DAcc-level1 AccDef)

qed

lemma A71-NSources-L1:

$\forall C \in (\text{AbstrLevel level1}). (C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA81 \longrightarrow sA71 \notin (\text{Sources level1 } C))$

proof –

have $\forall C \in (\text{AbstrLevel level1}). (C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA81 \longrightarrow (C \notin (\text{Acc level1 } sA71)))$

by (metis A71-Acc insert-iff singleton-iff)

thus ?thesis **by** (metis Acc-SourcesNOT)

qed

lemma A6-Acc-L1:

$(\text{Acc level1 } sA6) = \{sA71, sA72, sA81, sA82, sA91, sA92, sA93\}$

proof –

have daA6: $D\text{Acc level1 } sA6 = \{sA71, sA72\}$ **by** (rule A6-DAcc-level1)

hence $(\bigcup S \in (D\text{Acc level1 } sA6). (\text{Acc level1 } S)) = (\text{Acc level1 } sA71) \cup (\text{Acc level1 } sA72)$ **by** simp

hence aA6: $(\bigcup S \in (D\text{Acc level1 } sA6). (\text{Acc level1 } S)) = \{sA81, sA91, sA92\} \cup \{sA82, sA93\}$

by (simp add: A71-Acc A72-Acc)

have $(\text{Acc level1 } sA6) = (D\text{Acc level1 } sA6) \cup (\bigcup S \in (D\text{Acc level1 } sA6). (\text{Acc level1 } S))$

by (rule AccDef)

with daA6 aA6 **show** ?thesis **by** auto

qed

lemma *A6-NSources-L1Acc*:

$\forall C \in (\text{AbstrLevel } \text{level1}). (C \notin (\text{Acc } \text{level1 } sA6) \longrightarrow sA6 \notin (\text{Sources } \text{level1 } C))$
by (*metis Acc-SourcesNOT*)

lemma *A6-NSources-L1*:

$\forall C \in (\text{AbstrLevel } \text{level1}). (C \neq sA93 \wedge C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA82 \wedge C \neq sA81 \wedge C \neq sA72 \wedge C \neq sA71 \longrightarrow sA6 \notin (\text{Sources } \text{level1 } C))$

proof –

have $\forall C \in (\text{AbstrLevel } \text{level1}). (C \neq sA93 \wedge C \neq sA92 \wedge C \neq sA91 \wedge C \neq sA82 \wedge C \neq sA81 \wedge C \neq sA72 \wedge C \neq sA71$

$\longrightarrow (C \notin (\text{Acc } \text{level1 } sA6)))$

by (*metis A6-Acc-L1 empty-iff insert-iff*)

thus *?thesis* **by** (*metis Acc-SourcesNOT*)

qed

lemma *A5-Acc-L1*: $(\text{Acc } \text{level1 } sA5) = \{\}$

by (*metis A5-NotSource-level1 isNotDSource-EmptyAcc*)

lemma *SourcesA11-L1*: $\text{Sources } \text{level1 } sA11 = \{\}$

by (*simp add: DSourcesA11-L1 DSourcesEmptySources*)

lemma *SourcesA12-L1*: $\text{Sources } \text{level1 } sA12 = \{\}$

by (*simp add: DSourcesA12-L1 DSourcesEmptySources*)

lemma *SourcesA21-L1*: $\text{Sources } \text{level1 } sA21 = \{sA11\}$

by (*simp add: DSourcesA21-L1 SourcesA11-L1 Sources-singleDSource*)

lemma *SourcesA22-L1*: $\text{Sources } \text{level1 } sA22 = \{sA11, sA22, sA23, sA31, sA32, sA41\}$

proof

show $\text{Sources } \text{level1 } sA22 \subseteq \{sA11, sA22, sA23, sA31, sA32, sA41\}$

proof –

have $A2\text{level1}:sA22 \in (\text{AbstrLevel } \text{level1})$ **by** (*simp add: AbstrLevel1*)

from $A2\text{level1}$ **have** $sgA42:sA42 \notin \text{Sources } \text{level1 } sA22$ **by** (*metis A42-NSources-L1 CSet.distinct(347)*)

have $sgA5:sA5 \notin \text{Sources } \text{level1 } sA22$

by (*metis A5-NotSource-level1 Acc-Sources all-not-in-conv isNotDSource-EmptyAcc*)

have $sgA12:sA12 \notin \text{Sources } \text{level1 } sA22$ **by** (*metis A12-NotSource-level1 A2level1 isNotSource-Sources*)

have $sgA21:sA21 \notin \text{Sources } \text{level1 } sA22$

by (*metis A21-NotSource-level1 DAcc-DSourcesNOT NDSourceExistsDSource empty-iff isNotDSource-EmptyDAcc*)

from $A2\text{level1}$ **have** $sgA6:sA6 \notin \text{Sources } \text{level1 } sA22$ **by** (*simp add: A6-NSources-L1*)

from *A2level1* **have** $sgA71:sA71 \notin Sources\ level1\ sA22$ **by** (*simp add: A71-NSources-L1*)
from *A2level1* **have** $sgA72:sA72 \notin Sources\ level1\ sA22$ **by** (*simp add: A72-NSources-L1*)
from *A2level1* **have** $sgA81:sA81 \notin Sources\ level1\ sA22$ **by** (*simp add: A81-NSources-L1*)
from *A2level1* **have** $sgA82:sA82 \notin Sources\ level1\ sA22$ **by** (*simp add: A82-NSources-L1*)
from *A2level1* **have** $sgA91:sA91 \notin Sources\ level1\ sA22$ **by** (*simp add: A91-NSources-L1*)
from *A2level1* **have** $sgA92:sA92 \notin Sources\ level1\ sA22$ **by** (*simp add: A92-NSources-L1*)
from *A2level1* **have** $sgA93:sA93 \notin Sources\ level1\ sA22$ **by** (*metis A93-NotSourceSet-level1*)

have $Sources\ level1\ sA22 \subseteq \{sA11, sA12, sA21, sA22, sA23, sA31, sA32, sA41, sA42, sA5, sA6, sA71, sA72, sA81, sA82, sA91, sA92, sA93\}$
by (*metis AbstrLevel1 SourcesLevelX*)
with *sgA5 sgA12 sgA21 sgA42 sgA6 sgA71 sgA72 sgA81 sgA82 sgA91 sgA92 sgA93* **show**
 $Sources\ level1\ sA22 \subseteq \{sA11, sA22, sA23, sA31, sA32, sA41\}$
by *auto*

qed

next

show $\{sA11, sA22, sA23, sA31, sA32, sA41\} \subseteq Sources\ level1\ sA22$
proof –
have $sDef:(Sources\ level1\ sA22) = (DSources\ level1\ sA22) \cup (\bigcup S \in (DSources\ level1\ sA22). (Sources\ level1\ S))$
by (*rule SourcesDef*)
have *A11s*: $sA11 \in Sources\ level1\ sA22$ **by** (*metis DSourceIsSource DSourcesA22-L1 insertI1*)
have *A41s*: $sA41 \in Sources\ level1\ sA22$ **by** (*metis (full-types) DSourceIsSource DSourcesA22-L1 insertCI*)
have *A31s*: $sA31 \in Sources\ level1\ sA22$
by (*metis (full-types) A41s DSourceIsSource DSourcesA41-L1 SourcesTrans insertCI*)
have *A32s*: $sA32 \in Sources\ level1\ sA22$
by (*metis A32-DAcc-level1 A41s DAcc-DSourcesNOT DSourceOfSource insertI1*)
have *A23s*: $sA23 \in Sources\ level1\ sA22$ **by** (*metis A32s DSourceOfSource DSourcesA32-L1 insertI1*)
have *A22s*: $sA22 \in Sources\ level1\ sA22$ **by** (*metis A31s DSourceOfSource DSourcesA31-L1 insertI1*)
with *A11s A22s A23s A31s A32s A41s* **show** *?thesis* **by** *auto*
qed
qed

lemma *SourcesA23-L1*: $Sources\ level1\ sA23 = \{sA11\}$
by (*simp add: DSourcesA23-L1 SourcesA11-L1 Sources-singleDSource*)

lemma *SourcesA31-L1*: Sources level1 $sA31 = \{sA11, sA22, sA23, sA31, sA32, sA41\}$

by (*metis DSourcesA31-L1 SourcesA22-L1 Sources-singleDSource Un-insert-right insert-absorb2 insert-is-Un*)

lemma *SourcesA32-L1*: Sources level1 $sA32 = \{sA11, sA23\}$

by (*metis DSourcesA32-L1 SourcesA23-L1 Sources-singleDSource Un-insert-right insert-is-Un*)

lemma *SourcesA41-L1*: Sources level1 $sA41 = \{sA11, sA22, sA23, sA31, sA32, sA41\}$

by (*metis DSourcesA41-L1 SourcesA31-L1 SourcesA32-L1 Sources-2DSources Un-absorb Un-commute Un-insert-left*)

lemma *SourcesA42-L1*: Sources level1 $sA42 = \{\}$

by (*simp add: DSourcesA42-L1 DSourcesEmptySources*)

lemma *SourcesA5-L1*: Sources level1 $sA5 = \{sA42\}$

by (*simp add: DSourcesA5-L1 SourcesA42-L1 Sources-singleDSource*)

lemma *SourcesA6-L1*: Sources level1 $sA6 = \{\}$

by (*simp add: DSourcesA6-L1 DSourcesEmptySources*)

lemma *SourcesA71-L1*: Sources level1 $sA71 = \{sA6\}$

by (*metis DSourcesA71-L1 SourcesA6-L1 SourcesEmptyDSources SourcesOnlyD-Sources singleton-iff*)

lemma *SourcesA81-L1*: Sources level1 $sA81 = \{sA6, sA71, sA81, sA91\}$

proof –

have *dA81:DSources level1* $sA81 = \{sA71, sA91\}$ **by** (*rule DSourcesA81-L1*)

have *dA91:DSources level1* $sA91 = \{sA81\}$ **by** (*rule DSourcesA91-L1*)

have (*Sources level1* $sA81$) = (*DSources level1* $sA81$) \cup ($\bigcup S \in$ (*DSources level1* $sA81$). (*Sources level1* S))

by (*rule SourcesDef*)

with *dA81* **have** (*Sources level1* $sA81$) = ($\{sA71, sA91\} \cup$ (*Sources level1* $sA71$) \cup (*Sources level1* $sA91$))

by (*metis (hide-lams, no-types) SUP-empty UN-insert Un-insert-left sup-bot.left-neutral sup-commute*)

hence *sourcesA81*:(*Sources level1* $sA81$) = ($\{sA71, sA91, sA6\} \cup$ (*Sources level1* $sA91$))

by (*metis SourcesA71-L1 insert-is-Un sup-assoc*)

have (*Sources level1* $sA91$) = (*DSources level1* $sA91$) \cup ($\bigcup S \in$ (*DSources level1* $sA91$). (*Sources level1* S))

by (*rule SourcesDef*)

with *dA91* **have** (*Sources level1* $sA91$) = ($\{sA81\} \cup$ (*Sources level1* $sA81$)) **by** *simp*

with *sourcesA81* **have** (*Sources level1* $sA81$) = $\{sA71, sA91, sA6\} \cup \{sA81\} \cup \{sA81, sA91\}$

by (*metis SourcesLoop*)

thus *?thesis* **by** *auto*
qed

lemma *SourcesA91-L1: Sources level1 sA91 = {sA6, sA71, sA81, sA91}*

proof –

have *DSources level1 sA91 = {sA81}* **by** (*rule DSourcesA91-L1*)

thus *?thesis* **by** (*metis SourcesA81-L1 Sources-singleDSource*
Un-empty-left Un-insert-left insert-absorb2 insert-commute)

qed

lemma *SourcesA92-L1: Sources level1 sA92 = {sA6, sA71, sA81, sA91}*

by (*metis DSourcesA91-L1 DSourcesA92-L1 SourcesA91-L1 Sources-singleDSource*)

lemma *SourcesA72-L1: Sources level1 sA72 = {sA6}*

by (*metis DSourcesA6-L1 DSourcesA72-L1 SourcesOnlyDSources singleton-iff*)

lemma *SourcesA82-L1: Sources level1 sA82 = {sA6, sA72}*

proof –

have *dA82:DSources level1 sA82 = {sA72}* **by** (*rule DSourcesA82-L1*)

have (*Sources level1 sA82*) = (*DSources level1 sA82*) \cup ($\bigcup S \in$ (*DSources level1*
sA82). (*Sources level1 S*))

by (*rule SourcesDef*)

with *dA82* **have** (*Sources level1 sA82*) = {*sA72*} \cup (*Sources level1 sA72*) **by**
simp

thus *?thesis* **by** (*metis SourcesA72-L1 Un-commute insert-is-Un*)

qed

lemma *SourcesA93-L1: Sources level1 sA93 = {sA6, sA72, sA82}*

by (*metis DSourcesA93-L1 SourcesA82-L1 Sources-singleDSource Un-insert-right*
insert-is-Un)

— Abstraction level 2

lemma *SourcesS1-L2: Sources level2 sS1 = {}*

proof –

have *DSources level2 sS1 = {}* **by** (*simp add: DSources-def AbstrLevel2, auto*)

thus *?thesis* **by** (*simp add: DSourcesEmptySources*)

qed

lemma *SourcesS2-L2: Sources level2 sS2 = {}*

proof –

have *DSources level2 sS2 = {}* **by** (*simp add: DSources-def AbstrLevel2, auto*)

thus *?thesis* **by** (*simp add: DSourcesEmptySources*)

qed

lemma *SourcesS3-L2: Sources level2 sS3 = {sS2}*

proof –
have $DSourcesS3:DSources\ level2\ sS3 = \{sS2\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS2 = \{\}$ **by** (*rule SourcesS2-L2*)
with $DSourcesS3$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS4-L2$: $Sources\ level2\ sS4 = \{sS2\}$
proof –
have $DSourcesS4:DSources\ level2\ sS4 = \{sS2\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS2 = \{\}$ **by** (*rule SourcesS2-L2*)
with $DSourcesS4$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS5-L2$: $Sources\ level2\ sS5 = \{sS2, sS4\}$
proof –
have $DSourcesS5:DSources\ level2\ sS5 = \{sS4\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS4 = \{sS2\}$ **by** (*rule SourcesS4-L2*)
with $DSourcesS5$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS6-L2$: $Sources\ level2\ sS6 = \{sS2, sS4, sS5\}$
proof –
have $DSourcesS6:DSources\ level2\ sS6 = \{sS2, sS5\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $SourcesS2:Sources\ level2\ sS2 = \{\}$ **by** (*rule SourcesS2-L2*)
have $Sources\ level2\ sS5 = \{sS2, sS4\}$ **by** (*rule SourcesS5-L2*)
with $SourcesS2\ DSourcesS6$ **show** $?thesis$ **by** (*simp add: Sources-2DSources, auto*)
qed

lemma $SourcesS7-L2$: $Sources\ level2\ sS7 = \{\}$
proof –
have $DSources\ level2\ sS7 = \{\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
thus $?thesis$ **by** (*simp add: DSourcesEmptySources*)
qed

lemma $SourcesS8-L2$:
 $Sources\ level2\ sS8 = \{sS7\}$
proof –
have $DSourcesS8:DSources\ level2\ sS8 = \{sS7\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS7 = \{\}$ **by** (*rule SourcesS7-L2*)
with $DSourcesS8$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS9-L2$:

$Sources\ level2\ sS9 = \{\}$
proof –
have $DSources\ level2\ sS9 = \{\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
thus $?thesis$ **by** (*simp add: DSourcesEmptySources*)
qed

lemma $SourcesS10-L2$: $Sources\ level2\ sS10 = \{sS9\}$
proof –
have $DSourcesS10:DSources\ level2\ sS10 = \{sS9\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS9 = \{\}$ **by** (*rule SourcesS9-L2*)
with $DSourcesS10$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS11-L2$: $Sources\ level2\ sS11 = \{sS9\}$
proof –
have $DSourcesS11:DSources\ level2\ sS11 = \{sS9\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS9 = \{\}$ **by** (*rule SourcesS9-L2*)
with $DSourcesS11$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS12-L2$: $Sources\ level2\ sS12 = \{sS9, sS10\}$
proof –
have $DSourcesS12:DSources\ level2\ sS12 = \{sS10\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS10 = \{sS9\}$ **by** (*rule SourcesS10-L2*)
with $DSourcesS12$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS13-L2$: $Sources\ level2\ sS13 = \{sS9, sS10, sS12\}$
proof –
have $DSourcesS13:DSources\ level2\ sS13 = \{sS12\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS12 = \{sS9, sS10\}$ **by** (*rule SourcesS12-L2*)
with $DSourcesS13$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS14-L2$: $Sources\ level2\ sS14 = \{sS9, sS11\}$
proof –
have $DSourcesS14:DSources\ level2\ sS14 = \{sS11\}$ **by** (*simp add: DSources-def AbstrLevel2, auto*)
have $Sources\ level2\ sS11 = \{sS9\}$ **by** (*rule SourcesS11-L2*)
with $DSourcesS14$ **show** $?thesis$ **by** (*simp add: Sources-singleDSource*)
qed

lemma $SourcesS15-L2$: $Sources\ level2\ sS15 = \{sS9, sS11, sS14\}$
proof –
have $DSourcesS15:DSources\ level2\ sS15 = \{sS14\}$ **by** (*simp add: DSources-def*)

AbstrLevel2, auto)
have *Sources level2* $sS14 = \{sS9, sS11\}$ **by** (*rule SourcesS14-L2*)
with *DSourcesS15* **show** *?thesis* **by** (*simp add: Sources-singleDSource*)
qed

4.5 Minimal sets of components to prove certain properties

lemma *minSetOfComponentsTestL2p1*:
minSetOfComponents level2 $\{data10, data13\} = \{sS1\}$
proof –
have *outL2:outSetOfComponents level2* $\{data10, data13\} = \{sS1\}$
by (*simp add: outSetOfComponents-def AbstrLevel2, auto*)
have *Sources level2* $sS1 = \{\}$ **by** (*simp add: SourcesS1-L2*)
with *outL2* **show** *?thesis* **by** (*simp add: minSetOfComponents-def*)
qed

lemma *NOT-noIrrelevantChannelsTestL2p1*:
 \neg *noIrrelevantChannels level2* $\{data10, data13\}$
by (*simp add: noIrrelevantChannels-def systemIN-def minSetOfComponentsTestL2p1 AbstrLevel2*)

lemma *NOT-allNeededINChannelsTestL2p1*:
 \neg *allNeededINChannels level2* $\{data10, data13\}$
by (*simp add: allNeededINChannels-def minSetOfComponentsTestL2p1 systemIN-def AbstrLevel2*)

lemma *minSetOfComponentsTestL2p2*:
minSetOfComponents level2 $\{data1, data12\} = \{sS2, sS4, sS5, sS6\}$
proof –
have *outL2:outSetOfComponents level2* $\{data1, data12\} = \{sS6\}$
by (*simp add: outSetOfComponents-def AbstrLevel2, auto*)
have *Sources level2* $sS6 = \{sS2, sS4, sS5\}$
by (*simp add: SourcesS6-L2*)
with *outL2* **show** *?thesis*
by (*simp add: minSetOfComponents-def*)
qed

lemma *noIrrelevantChannelsTestL2p2*:
noIrrelevantChannels level2 $\{data1, data12\}$
by (*simp add: noIrrelevantChannels-def systemIN-def minSetOfComponentsTestL2p2 AbstrLevel2*)

lemma *allNeededINChannelsTestL2p2*:
allNeededINChannels level2 $\{data1, data12\}$
by (*simp add: allNeededINChannels-def minSetOfComponentsTestL2p2 systemIN-def AbstrLevel2*)

lemma *minSetOfComponentsTestL1p3*:
minSetOfComponents level1 $\{data1, data10, data11\} = \{sA12, sA11, sA21\}$

proof –
have $sg1.outSetOfComponents\ level1\ \{data1, data10, data11\} = \{sA12, sA21\}$
by (*simp add: outSetOfComponents-def AbstrLevel1, auto*)
have $DSources\ level1\ sA12 = \{\}$
by (*simp add: DSources-def AbstrLevel1, auto*)
hence $sg2.Sources\ level1\ sA12 = \{\}$
by (*simp add: DSourcesEmptySources*)
have $sg3.DSources\ level1\ sA21 = \{sA11\}$
by (*simp add: DSources-def AbstrLevel1, auto*)
have $sg4.DSources\ level1\ sA11 = \{\}$
by (*simp add: DSources-def AbstrLevel1, auto*)
hence $Sources\ level1\ sA21 = \{sA11\}$
by (*metis SourcesOnlyDSources sg3 singleton-iff*)
from this and sg1 and sg2 show *?thesis*
by (*simp add: minSetOfComponents-def, blast*)
qed

lemma *noIrrelevantChannelsTestL1p3*:
noIrrelevantChannels level1 {data1, data10, data11}
by (*simp add: noIrrelevantChannels-def systemIN-def minSetOfComponentsTestL1p3 AbstrLevel1*)

lemma *allNeededINChannelsTestL1p3*:
allNeededINChannels level1 {data1, data10, data11}
by (*simp add: allNeededINChannels-def minSetOfComponentsTestL1p3 systemIN-def AbstrLevel1*)

lemma *minSetOfComponentsTestL2p3*:
minSetOfComponents level2 {data1, data10, data11} = {sS1, sS2, sS3}

proof –
have $sg1.outSetOfComponents\ level2\ \{data1, data10, data11\} = \{sS1, sS3\}$
by (*simp add: outSetOfComponents-def AbstrLevel2, auto*)
have $sS1.Sources\ level2\ sS1 = \{\}$ **by** (*simp add: SourcesS1-L2*)
have $Sources\ level2\ sS3 = \{sS2\}$ **by** (*simp add: SourcesS3-L2*)
with $sg1\ sS1$ **show** *?thesis*
by (*simp add: minSetOfComponents-def, blast*)
qed

lemma *noIrrelevantChannelsTestL2p3*:
noIrrelevantChannels level2 {data1, data10, data11}
by (*simp add: noIrrelevantChannels-def systemIN-def minSetOfComponentsTestL2p3 AbstrLevel2*)

lemma *allNeededINChannelsTestL2p3*:
allNeededINChannels level2 {data1, data10, data11}
by (*simp add: allNeededINChannels-def minSetOfComponentsTestL2p3 systemIN-def AbstrLevel2*)

end

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